Legal Notice

PostgreSQL is Copyright © 1996-2018 by the PostgreSQL Global Development Group.

Postgres95 is Copyright © 1994-5 by the Regents of the University of California.

Permission to use, copy, modify, and distribute this software and its documentation for any purpose, without fee, and without a written agreement is hereby granted, provided that the above copyright notice and this paragraph and the following two paragraphs appear in all copies.

IN NO EVENT SHALL THE UNIVERSITY OF CALIFORNIA BE LIABLE TO ANY PARTY FOR DIRECT, INDIRECT, SPECIAL, INCIDENTAL, OR CONSEQUENTIAL DAMAGES, INCLUDING LOST PROFITS, ARISING OUT OF THE USE OF THIS SOFTWARE AND ITS DOCUMENTATION, EVEN IF THE UNIVERSITY OF CALIFORNIA HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGE.

THE UNIVERSITY OF CALIFORNIA SPECIFICALLY DISCLAIMS ANY WARRANTIES, INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE. THE SOFTWARE PROVIDED HEREUNDER IS ON AN "AS-IS" BASIS, AND THE UNIVERSITY OF CALIFORNIA HAS NO OBLIGATIONS TO PROVIDE MAINTENANCE, SUPPORT, UPDATES, ENHANCEMENTS, OR MODIFICATIONS.
# Table of Contents

Preface .................................................................................................................. xxxvi

1. What is PostgreSQL? ....................................................................................... xxxvi

2. A Brief History of PostgreSQL ..................................................................... xxxvi

   2.1. The Berkeley POSTGRES Project ......................................................... xxvii
   2.2. Postgres95 ............................................................................................. xxxii
   2.3. PostgreSQL ............................................................................................ xxviii

3. Conventions ...................................................................................................... xxviii

4. Bug Reporting Guidelines .................................................................................. xxxix

   5.1. Identifying Bugs ....................................................................................... xxxix
   5.2. What to Report ....................................................................................... xli
   5.3. Where to Report Bugs ........................................................................... xlii

I. Tutorial ............................................................................................................. 1

   1. Getting Started ........................................................................................... 3
       1.1. Installation ............................................................................................ 3
       1.2. Architectural Fundamentals ............................................................... 3
       1.3. Creating a Database ........................................................................... 3
       1.4. Accessing a Database ....................................................................... 5

   2. The SQL Language ....................................................................................... 7
       2.1. Introduction ....................................................................................... 7
       2.2. Concepts ............................................................................................ 7
       2.3. Creating a New Table ........................................................................ 7
       2.4. Populating a Table With Rows ......................................................... 8
       2.5. Querying a Table .............................................................................. 9
       2.6. Joins Between Tables ....................................................................... 11
       2.7. Aggregate Functions ........................................................................ 13
       2.8. Updates ............................................................................................. 15
       2.9. Deletions ............................................................................................ 15

   3. Advanced Features ....................................................................................... 16
       3.1. Introduction ....................................................................................... 16
       3.2. Views ................................................................................................. 16
       3.3. Foreign Keys ...................................................................................... 16
       3.4. Transactions ...................................................................................... 17
       3.5. Window Functions ............................................................................ 19
       3.6. Inheritance ....................................................................................... 22
       3.7. Conclusion ....................................................................................... 23

II. The SQL Language .......................................................................................... 24

   4. SQL Syntax .................................................................................................. 32
       4.1. Lexical Structure .............................................................................. 32
       4.2. Value Expressions ............................................................................ 41
       4.3. Calling Functions .............................................................................. 55

   5. Data Definition ............................................................................................ 58
       5.1. Table Basics ..................................................................................... 58
       5.2. Default Values .................................................................................. 59
       5.3. Constraints ....................................................................................... 60
       5.4. System Columns .............................................................................. 67
       5.5. Modifying Tables ............................................................................. 68
       5.6. Privileges ......................................................................................... 71
       5.7. Row Security Policies ..................................................................... 72
       5.8. Schemas ........................................................................................... 78
       5.9. Inheritance ....................................................................................... 83
       5.10. Table Partitioning .......................................................................... 86
       5.11. Foreign Data .................................................................................... 98
       5.12. Other Database Objects ................................................................. 99
       5.13. Dependency Tracking .................................................................... 99
6. Data Manipulation ................................................................. 101
   6.1. Inserting Data ............................................................. 101
   6.2. Updating Data ............................................................ 102
   6.3. Deleting Data ............................................................. 103
   6.4. Returning Data From Modified Rows ................................. 103

7. Queries ................................................................. 105
   7.1. Overview ................................................................. 105
   7.2. Table Expressions ...................................................... 105
   7.3. Select Lists ............................................................. 120
   7.4. Combining Queries ..................................................... 122
   7.5. Sorting Rows ............................................................. 122
   7.6. LIMIT and OFFSET ...................................................... 123
   7.7. VALUES Lists ............................................................ 124
   7.8. WITH Queries (Common Table Expressions) .......................... 125

8. Data Types .......................................................... 131
   8.1. Numeric Types .......................................................... 132
   8.2. Monetary Types ......................................................... 137
   8.3. Character Types ......................................................... 138
   8.4. Binary Data Types ...................................................... 140
   8.5. Date/Time Types ......................................................... 142
   8.6. Boolean Type ............................................................. 151
   8.7. Enumerated Types ....................................................... 152
   8.8. Geometric Types ......................................................... 154
   8.9. Network Address Types ................................................ 156
   8.10. Bit String Types ......................................................... 159
   8.11. Text Search Types ..................................................... 160
   8.12. UUID Type ............................................................... 162
   8.13. XML Type ............................................................... 163
   8.14. JSON Types ............................................................. 165
   8.15. Arrays ................................................................. 172
   8.16. Composite Types ....................................................... 181
   8.17. Range Types ............................................................ 188
   8.18. Domain Types .......................................................... 193
   8.19. Object Identifier Types ............................................... 194
   8.20. pg_lsn Type .............................................................. 196
   8.21. Pseudo-Types .......................................................... 196

9. Functions and Operators ................................................. 198
   9.1. Logical Operators ....................................................... 198
   9.2. Comparison Functions and Operators ................................ 198
   9.3. Mathematical Functions and Operators ................................ 201
   9.4. String Functions and Operators ...................................... 204
   9.5. Binary String Functions and Operators ............................. 219
   9.6. Bit String Functions and Operators ................................ 221
   9.7. Pattern Matching ....................................................... 222
   9.8. Data Type Formatting Functions ..................................... 237
   9.9. Date/Time Functions and Operators ................................ 244
   9.10. Enum Support Functions ............................................... 257
   9.11. Geometric Functions and Operators ................................ 258
   9.12. Network Address Functions and Operators ......................... 262
   9.13. Text Search Functions and Operators ............................. 264
   9.14. XML Functions ......................................................... 271
   9.15. JSON Functions and Operators ..................................... 284
   9.16. Sequence Manipulation Functions .................................. 293
   9.17. Conditional Expressions .............................................. 295
   9.18. Array Functions and Operators ..................................... 298
   9.19. Range Functions and Operators .................................... 301
   9.20. Aggregate Functions ................................................ 303
   9.21. Window Functions .................................................... 310
# PostgreSQL 11.1 Documentation

## Table of Contents

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>9.22. Subquery Expressions</td>
<td>312</td>
</tr>
<tr>
<td>9.23. Row and Array Comparisons</td>
<td>315</td>
</tr>
<tr>
<td>9.24. Set Returning Functions</td>
<td>318</td>
</tr>
<tr>
<td>9.25. System Information Functions</td>
<td>321</td>
</tr>
<tr>
<td>9.26. System Administration Functions</td>
<td>338</td>
</tr>
<tr>
<td>9.27. Trigger Functions</td>
<td>355</td>
</tr>
<tr>
<td>9.28. Event Trigger Functions</td>
<td>355</td>
</tr>
<tr>
<td>10. Type Conversion</td>
<td>359</td>
</tr>
<tr>
<td>10.1. Overview</td>
<td>359</td>
</tr>
<tr>
<td>10.2. Operators</td>
<td>360</td>
</tr>
<tr>
<td>10.3. Functions</td>
<td>364</td>
</tr>
<tr>
<td>10.4. Value Storage</td>
<td>368</td>
</tr>
<tr>
<td>10.5. \texttt{UNION}, \texttt{CASE}, and Related Constructs</td>
<td>368</td>
</tr>
<tr>
<td>10.6. \texttt{SELECT} Output Columns</td>
<td>370</td>
</tr>
<tr>
<td>11. Indexes</td>
<td>371</td>
</tr>
<tr>
<td>11.1. Introduction</td>
<td>371</td>
</tr>
<tr>
<td>11.2. Index Types</td>
<td>372</td>
</tr>
<tr>
<td>11.3. Multicolumn Indexes</td>
<td>374</td>
</tr>
<tr>
<td>11.4. Indexes and ORDER BY</td>
<td>375</td>
</tr>
<tr>
<td>11.5. Combining Multiple Indexes</td>
<td>375</td>
</tr>
<tr>
<td>11.6. Unique Indexes</td>
<td>376</td>
</tr>
<tr>
<td>11.7. Indexes on Expressions</td>
<td>377</td>
</tr>
<tr>
<td>11.8. Partial Indexes</td>
<td>377</td>
</tr>
<tr>
<td>11.9. Operator Classes and Operator Families</td>
<td>380</td>
</tr>
<tr>
<td>11.10. Indexes and Collations</td>
<td>381</td>
</tr>
<tr>
<td>11.11. Index-Only Scans</td>
<td>382</td>
</tr>
<tr>
<td>11.12. Examining Index Usage</td>
<td>384</td>
</tr>
<tr>
<td>12. Full Text Search</td>
<td>385</td>
</tr>
<tr>
<td>12.1. Introduction</td>
<td>385</td>
</tr>
<tr>
<td>12.2. Tables and Indexes</td>
<td>389</td>
</tr>
<tr>
<td>12.3. Controlling Text Search</td>
<td>391</td>
</tr>
<tr>
<td>12.4. Additional Features</td>
<td>398</td>
</tr>
<tr>
<td>12.5. Parsers</td>
<td>403</td>
</tr>
<tr>
<td>12.6. Dictionaries</td>
<td>405</td>
</tr>
<tr>
<td>12.7. Configuration Example</td>
<td>415</td>
</tr>
<tr>
<td>12.8. Testing and Debugging Text Search</td>
<td>416</td>
</tr>
<tr>
<td>12.9. GIN and GiST Index Types</td>
<td>421</td>
</tr>
<tr>
<td>12.10. \texttt{psql} Support</td>
<td>421</td>
</tr>
<tr>
<td>12.11. Limitations</td>
<td>424</td>
</tr>
<tr>
<td>13. Concurrency Control</td>
<td>426</td>
</tr>
<tr>
<td>13.1. Introduction</td>
<td>426</td>
</tr>
<tr>
<td>13.2. Transaction Isolation</td>
<td>426</td>
</tr>
<tr>
<td>13.3. Explicit Locking</td>
<td>432</td>
</tr>
<tr>
<td>13.4. Data Consistency Checks at the Application Level</td>
<td>437</td>
</tr>
<tr>
<td>13.5. Caveats</td>
<td>439</td>
</tr>
<tr>
<td>13.6. Locking and Indexes</td>
<td>439</td>
</tr>
<tr>
<td>14. Performance Tips</td>
<td>441</td>
</tr>
<tr>
<td>14.1. Using \texttt{EXPLAIN}</td>
<td>441</td>
</tr>
<tr>
<td>14.2. Statistics Used by the Planner</td>
<td>452</td>
</tr>
<tr>
<td>14.3. Controlling the Planner with Explicit JOIN Clauses</td>
<td>456</td>
</tr>
<tr>
<td>14.4. Populating a Database</td>
<td>458</td>
</tr>
<tr>
<td>14.5. Non-Durable Settings</td>
<td>461</td>
</tr>
<tr>
<td>15. Parallel Query</td>
<td>462</td>
</tr>
<tr>
<td>15.1. How Parallel Query Works</td>
<td>462</td>
</tr>
<tr>
<td>15.2. When Can Parallel Query Be Used?</td>
<td>463</td>
</tr>
<tr>
<td>15.3. Parallel Plans</td>
<td>464</td>
</tr>
<tr>
<td>15.4. Parallel Safety</td>
<td>466</td>
</tr>
<tr>
<td>III. Server Administration</td>
<td>468</td>
</tr>
</tbody>
</table>
16. Installation from Source Code .......................................................... 474
16.1. Short Version ........................................................................... 474
16.2. Requirements ........................................................................ 474
16.3. Getting The Source ................................................................. 476
16.4. Installation Procedure ............................................................ 476
16.5. Post-Installation Setup ............................................................ 488
16.6. Supported Platforms ............................................................... 490
16.7. Platform-specific Notes .......................................................... 490
17. Installation from Source Code on Windows ................................. 498
17.1. Building with Visual C++ or the Microsoft Windows SDK .......... 498
18. Server Setup and Operation .......................................................... 504
18.1. The PostgreSQL User Account ................................................. 504
18.2. Creating a Database Cluster .................................................... 504
18.3. Starting the Database Server .................................................. 506
18.4. Managing Kernel Resources .................................................. 509
18.5. Shutting Down the Server ...................................................... 518
18.6. Upgrading a PostgreSQL Cluster ........................................... 519
18.7. Preventing Server Spoofing .................................................... 522
18.8. Encryption Options ............................................................... 522
18.9. Secure TCP/IP Connections with SSL ..................................... 523
18.10. Secure TCP/IP Connections with SSH Tunnels ....................... 527
18.11. Registering Event Log on Windows ........................................ 528
19. Server Configuration ................................................................. 529
19.1. Setting Parameters ............................................................... 529
19.2. File Locations ...................................................................... 532
19.3. Connections and Authentication ........................................... 533
19.4. Resource Consumption ........................................................ 539
19.5. Write Ahead Log .................................................................. 545
19.6. Replication ............................................................................ 551
19.7. Query Planning ..................................................................... 556
19.8. Error Reporting and Logging ................................................ 562
19.9. Run-time Statistics ............................................................... 572
19.10. Automatic Vacuuming ......................................................... 574
19.11. Client Connection Defaults .................................................. 575
19.12. Lock Management ............................................................... 584
19.13. Version and Platform Compatibility ..................................... 585
19.15. Preset Options .................................................................... 587
19.16. Customized Options ............................................................ 589
19.17. Developer Options ............................................................. 589
19.18. Short Options ..................................................................... 592
20. Client Authentication ................................................................. 594
20.1. The pg_hba.conf File ............................................................. 594
20.2. User Name Maps ................................................................. 601
20.3. Authentication Methods ....................................................... 602
20.4. Trust Authentication ............................................................ 603
20.5. Password Authentication ...................................................... 603
20.6. GSSAPI Authentication ......................................................... 604
20.7. SSPI Authentication .............................................................. 605
20.8. Ident Authentication .............................................................. 606
20.9. Peer Authentication .............................................................. 607
20.10. LDAP Authentication ......................................................... 607
20.11. RADIUS Authentication ...................................................... 610
20.13. PAM Authentication ........................................................... 611
20.14. BSD Authentication ........................................................... 612
20.15. Authentication Problems .................................................... 612
21. Database Roles ........................................................................ 614
<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>31.8. Configuration Settings</td>
<td>752</td>
</tr>
<tr>
<td>31.9. Quick Setup</td>
<td>752</td>
</tr>
<tr>
<td>32. Just-in-Time Compilation (JIT)</td>
<td>754</td>
</tr>
<tr>
<td>32.1. What is JIT compilation?</td>
<td>754</td>
</tr>
<tr>
<td>32.2. When to JIT?</td>
<td>754</td>
</tr>
<tr>
<td>32.3. Configuration</td>
<td>756</td>
</tr>
<tr>
<td>32.4. Extensibility</td>
<td>756</td>
</tr>
<tr>
<td>33. Regression Tests</td>
<td>757</td>
</tr>
<tr>
<td>33.1. Running the Tests</td>
<td>757</td>
</tr>
<tr>
<td>33.2. Test Evaluation</td>
<td>760</td>
</tr>
<tr>
<td>33.3. Variant Comparison Files</td>
<td>763</td>
</tr>
<tr>
<td>33.4. TAP Tests</td>
<td>764</td>
</tr>
<tr>
<td>33.5. Test Coverage Examination</td>
<td>764</td>
</tr>
<tr>
<td>IV. Client Interfaces</td>
<td>765</td>
</tr>
<tr>
<td>34. libpq - C Library</td>
<td>770</td>
</tr>
<tr>
<td>34.1. Database Connection Control Functions</td>
<td>770</td>
</tr>
<tr>
<td>34.2. Connection Status Functions</td>
<td>783</td>
</tr>
<tr>
<td>34.3. Command Execution Functions</td>
<td>789</td>
</tr>
<tr>
<td>34.4. Asynchronous Command Processing</td>
<td>805</td>
</tr>
<tr>
<td>34.5. Retrieving Query Results Row-By-Row</td>
<td>809</td>
</tr>
<tr>
<td>34.6. Canceling Queries in Progress</td>
<td>810</td>
</tr>
<tr>
<td>34.7. The Fast-Path Interface</td>
<td>811</td>
</tr>
<tr>
<td>34.8. Asynchronous Notification</td>
<td>812</td>
</tr>
<tr>
<td>34.9. Functions Associated with the COPY Command</td>
<td>813</td>
</tr>
<tr>
<td>34.10. Control Functions</td>
<td>817</td>
</tr>
<tr>
<td>34.11. Miscellaneous Functions</td>
<td>819</td>
</tr>
<tr>
<td>34.12. Notice Processing</td>
<td>822</td>
</tr>
<tr>
<td>34.13. Event System</td>
<td>823</td>
</tr>
<tr>
<td>34.14. Environment Variables</td>
<td>829</td>
</tr>
<tr>
<td>34.15. The Password File</td>
<td>831</td>
</tr>
<tr>
<td>34.16. The Connection Service File</td>
<td>831</td>
</tr>
<tr>
<td>34.17. LDAP Lookup of Connection Parameters</td>
<td>832</td>
</tr>
<tr>
<td>34.18. SSL Support</td>
<td>833</td>
</tr>
<tr>
<td>34.19. Behavior in Threaded Programs</td>
<td>836</td>
</tr>
<tr>
<td>34.20. Building libpq Programs</td>
<td>837</td>
</tr>
<tr>
<td>34.21. Example Programs</td>
<td>838</td>
</tr>
<tr>
<td>35. Large Objects</td>
<td>850</td>
</tr>
<tr>
<td>35.1. Introduction</td>
<td>850</td>
</tr>
<tr>
<td>35.2. Implementation Features</td>
<td>850</td>
</tr>
<tr>
<td>35.3. Client Interfaces</td>
<td>850</td>
</tr>
<tr>
<td>35.4. Server-side Functions</td>
<td>854</td>
</tr>
<tr>
<td>35.5. Example Program</td>
<td>856</td>
</tr>
<tr>
<td>36. ECPG - Embedded SQL in C</td>
<td>862</td>
</tr>
<tr>
<td>36.1. The Concept</td>
<td>862</td>
</tr>
<tr>
<td>36.2. Managing Database Connections</td>
<td>862</td>
</tr>
<tr>
<td>36.3. Running SQL Commands</td>
<td>865</td>
</tr>
<tr>
<td>36.4. Using Host Variables</td>
<td>868</td>
</tr>
<tr>
<td>36.5. Dynamic SQL</td>
<td>882</td>
</tr>
<tr>
<td>36.6. pgtypes Library</td>
<td>884</td>
</tr>
<tr>
<td>36.7. Using Descriptor Areas</td>
<td>898</td>
</tr>
<tr>
<td>36.8. Error Handling</td>
<td>911</td>
</tr>
<tr>
<td>36.9. Preprocessor Directives</td>
<td>918</td>
</tr>
<tr>
<td>36.10. Processing Embedded SQL Programs</td>
<td>920</td>
</tr>
<tr>
<td>36.11. Library Functions</td>
<td>921</td>
</tr>
<tr>
<td>36.12. Large Objects</td>
<td>922</td>
</tr>
<tr>
<td>36.13. C++ Applications</td>
<td>923</td>
</tr>
<tr>
<td>36.14. Embedded SQL Commands</td>
<td>927</td>
</tr>
<tr>
<td>36.15. Informix Compatibility Mode</td>
<td>951</td>
</tr>
<tr>
<td>Section</td>
<td>Page</td>
</tr>
<tr>
<td>---------</td>
<td>------</td>
</tr>
<tr>
<td>36.16. Internals</td>
<td>966</td>
</tr>
<tr>
<td>37. The Information Schema</td>
<td>969</td>
</tr>
<tr>
<td>37.1. The Schema</td>
<td>969</td>
</tr>
<tr>
<td>37.2. Data Types</td>
<td>969</td>
</tr>
<tr>
<td>37.3. information_schema_catalog_name</td>
<td>970</td>
</tr>
<tr>
<td>37.4. administrable_role_authorizations</td>
<td>970</td>
</tr>
<tr>
<td>37.5. applicable_roles</td>
<td>970</td>
</tr>
<tr>
<td>37.6. attributes</td>
<td>971</td>
</tr>
<tr>
<td>37.7. character_sets</td>
<td>974</td>
</tr>
<tr>
<td>37.8. check_constraint_routine_usage</td>
<td>975</td>
</tr>
<tr>
<td>37.9. check_constraints</td>
<td>975</td>
</tr>
<tr>
<td>37.10. collations</td>
<td>976</td>
</tr>
<tr>
<td>37.11. collation_character_set_applicability</td>
<td>976</td>
</tr>
<tr>
<td>37.12. column_domain_usage</td>
<td>977</td>
</tr>
<tr>
<td>37.13. column_options</td>
<td>977</td>
</tr>
<tr>
<td>37.14. column_privileges</td>
<td>977</td>
</tr>
<tr>
<td>37.15. column_udt_usage</td>
<td>978</td>
</tr>
<tr>
<td>37.16. columns</td>
<td>979</td>
</tr>
<tr>
<td>37.17. constraint_column_usage</td>
<td>983</td>
</tr>
<tr>
<td>37.18. constraint_table_usage</td>
<td>983</td>
</tr>
<tr>
<td>37.19. data_type_privileges</td>
<td>984</td>
</tr>
<tr>
<td>37.20. domain_constraints</td>
<td>985</td>
</tr>
<tr>
<td>37.21. domain_udt_usage</td>
<td>985</td>
</tr>
<tr>
<td>37.22. domains</td>
<td>986</td>
</tr>
<tr>
<td>37.23. element_types</td>
<td>988</td>
</tr>
<tr>
<td>37.24. enabled_roles</td>
<td>991</td>
</tr>
<tr>
<td>37.25. foreign_data_wrapper_options</td>
<td>991</td>
</tr>
<tr>
<td>37.26. foreign_data_wrappers</td>
<td>991</td>
</tr>
<tr>
<td>37.27. foreign_server_options</td>
<td>992</td>
</tr>
<tr>
<td>37.28. foreign_servers</td>
<td>992</td>
</tr>
<tr>
<td>37.29. foreign_table_options</td>
<td>993</td>
</tr>
<tr>
<td>37.30. foreign_tables</td>
<td>993</td>
</tr>
<tr>
<td>37.31. key_column_usage</td>
<td>993</td>
</tr>
<tr>
<td>37.32. parameters</td>
<td>994</td>
</tr>
<tr>
<td>37.33. referential_constraints</td>
<td>996</td>
</tr>
<tr>
<td>37.34. role_column_grants</td>
<td>997</td>
</tr>
<tr>
<td>37.35. role_routine_grants</td>
<td>998</td>
</tr>
<tr>
<td>37.36. role_table_grants</td>
<td>998</td>
</tr>
<tr>
<td>37.37. role_udt_grants</td>
<td>999</td>
</tr>
<tr>
<td>37.38. role_usage_grants</td>
<td>1000</td>
</tr>
<tr>
<td>37.39. routine_privileges</td>
<td>1000</td>
</tr>
<tr>
<td>37.40. routines</td>
<td>1001</td>
</tr>
<tr>
<td>37.41. schemata</td>
<td>1006</td>
</tr>
<tr>
<td>37.42. sequences</td>
<td>1006</td>
</tr>
<tr>
<td>37.43. sql_features</td>
<td>1007</td>
</tr>
<tr>
<td>37.44. sql_implementation_info</td>
<td>1008</td>
</tr>
<tr>
<td>37.45. sql_languages</td>
<td>1008</td>
</tr>
<tr>
<td>37.46. sql_packages</td>
<td>1009</td>
</tr>
<tr>
<td>37.47. sql_parts</td>
<td>1009</td>
</tr>
<tr>
<td>37.48. sql_sizing</td>
<td>1010</td>
</tr>
<tr>
<td>37.49. sql-sizing_profiles</td>
<td>1010</td>
</tr>
<tr>
<td>37.50. table_constraints</td>
<td>1011</td>
</tr>
<tr>
<td>37.51. table_privileges</td>
<td>1011</td>
</tr>
<tr>
<td>37.52. tables</td>
<td>1012</td>
</tr>
<tr>
<td>37.53. transforms</td>
<td>1013</td>
</tr>
<tr>
<td>37.54. triggered_update_columns</td>
<td>1014</td>
</tr>
<tr>
<td>37.55. triggers</td>
<td>1014</td>
</tr>
<tr>
<td>37.56. udt_privileges</td>
<td>1016</td>
</tr>
</tbody>
</table>
V. Server Programming .......................................................... 1023
38. Extending SQL .............................................................. 1028
  38.1. How Extensibility Works ............................................. 1028
  38.2. The PostgreSQL Type System ...................................... 1028
  38.3. User-defined Functions .............................................. 1030
  38.4. User-defined Procedures ........................................... 1030
  38.5. Query Language (SQL) Functions .................................. 1030
  38.6. Function Overloading ............................................... 1046
  38.7. Function Volatility Categories .................................... 1047
  38.8. Procedural Language Functions .................................... 1049
  38.9. Internal Functions .................................................. 1049
  38.10. C-Language Functions .............................................. 1049
  38.11. User-defined Aggregates ........................................... 1070
  38.12. User-defined Types ................................................ 1077
  38.13. User-defined Operators ........................................... 1081
  38.14. Operator Optimization Information ............................... 1082
  38.15. Interfacing Extensions To Indexes ................................. 1086
  38.16. Packaging Related Objects into an Extension .................... 1099
  38.17. Extension Building Infrastructure ................................. 1106
39. Triggers ................................................................. 1110
  39.1. Overview of Trigger Behavior ..................................... 1110
  39.2. Visibility of Data Changes ........................................ 1113
  39.3. Writing Trigger Functions in C .................................... 1113
  39.4. A Complete Trigger Example ...................................... 1116
40. Event Triggers ............................................................. 1120
  40.1. Overview of Event Trigger Behavior ............................... 1120
  40.2. Event Trigger Firing Matrix ....................................... 1121
  40.3. Writing Event Trigger Functions in C ............................ 1125
  40.4. A Complete Event Trigger Example ............................... 1126
  40.5. A Table Rewrite Event Trigger Example ......................... 1127
41. The Rule System .......................................................... 1129
  41.1. The Query Tree ...................................................... 1129
  41.2. Views and the Rule System ........................................ 1130
  41.3. Materialized Views ................................................ 1137
  41.4. Rules on INSERT, UPDATE, and DELETE ......................... 1140
  41.5. Rules and Privileges ............................................... 1151
  41.6. Rules and Command Status ........................................ 1153
  41.7. Rules Versus Triggers ............................................... 1153
42. Procedural Languages .................................................. 1156
  42.1. Installing Procedural Languages .................................. 1156
43. PL/pgSQL - SQL Procedural Language .................................. 1159
  43.1. Overview ............................................................ 1159
  43.2. Structure of PL/pgSQL .............................................. 1160
  43.3. Declarations ....................................................... 1162
  43.4. Expressions ........................................................ 1167
  43.5. Basic Statements .................................................. 1168
  43.6. Control Structures ................................................ 1176
  43.7. Cursors ............................................................. 1190
  43.8. Transaction Management ........................................... 1196
  43.9. Errors and Messages ............................................... 1197
43.10. Trigger Functions ................................................................. 1199
43.11. PL/pgSQL Under the Hood .................................................. 1208
43.12. Tips for Developing in PL/pgSQL ........................................ 1211
43.13. Porting from Oracle PL/SQL .................................................. 1214
44. PL/Tcl - Tcl Procedural Language ............................................. 1224
  44.1. Overview ........................................................................ 1224
  44.2. PL/Tcl Functions and Arguments ........................................ 1224
  44.3. Data Values in PL/Tcl ....................................................... 1226
  44.4. Global Data in PL/Tcl ....................................................... 1226
  44.5. Database Access from PL/Tcl ............................................ 1227
  44.6. Trigger Functions in PL/Tcl .............................................. 1229
  44.7. Event Trigger Functions in PL/Tcl ..................................... 1231
  44.8. Error Handling in PL/Tcl .................................................. 1231
  44.9. Explicit Subtransactions in PL/Tcl ..................................... 1232
  44.10. Transaction Management ................................................ 1233
  44.11. PL/Tcl Configuration ..................................................... 1234
  44.12. Tcl Procedure Names ..................................................... 1234
45. PL/Perl - Perl Procedural Language ........................................... 1235
  45.1. PL/Perl Functions and Arguments ...................................... 1235
  45.2. Data Values in PL/Perl ..................................................... 1239
  45.3. Built-in Functions .......................................................... 1239
  45.4. Global Values in PL/Perl .................................................. 1244
  45.5. Trusted and Untrusted PL/Perl ........................................... 1245
  45.6. PL/Perl Triggers ............................................................. 1246
  45.7. PL/Perl Event Triggers .................................................... 1248
  45.8. PL/Perl Under the Hood .................................................. 1248
46. PL/Python - Python Procedural Language .................................... 1251
  46.1. Python 2 vs. Python 3 ...................................................... 1251
  46.2. PL/Python Functions ...................................................... 1252
  46.3. Data Values ................................................................. 1253
  46.4. Sharing Data ................................................................... 1259
  46.5. Anonymous Code Blocks ................................................ 1259
  46.6. Trigger Functions ........................................................... 1259
  46.7. Database Access ............................................................ 1260
  46.8. Explicit Subtransactions .................................................. 1264
  46.9. Transaction Management ................................................ 1265
  46.10. Utility Functions ............................................................ 1266
  46.11. Environment Variables .................................................. 1267
47. Server Programming Interface .................................................... 1269
  47.1. Interface Functions ......................................................... 1269
  47.2. Interface Support Functions ............................................. 1303
  47.3. Memory Management ...................................................... 1312
  47.4. Transaction Management ................................................ 1322
  47.5. Visibility of Data Changes ............................................... 1325
  47.6. Examples ................................................................. 1325
48. Background Worker Processes .................................................... 1329
49. Logical Decoding ................................................................. 1332
  49.1. Logical Decoding Examples ............................................... 1332
  49.2. Logical Decoding Concepts ............................................... 1334
  49.3. Streaming Replication Protocol Interface ............................ 1336
  49.4. Logical Decoding SQL Interface ........................................ 1336
  49.5. System Catalogs Related to Logical Decoding ....................... 1336
  49.6. Logical Decoding Output Plugins ..................................... 1336
  49.7. Logical Decoding Output Writers ..................................... 1341
  49.8. Synchronous Replication Support for Logical Decoding .......... 1341
50. Replication Progress Tracking .................................................. 1342
VI. Reference .............................................................................. 1343
  I. SQL Commands .................................................................... 1348
<table>
<thead>
<tr>
<th>Command</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>ABORT</td>
<td>1352</td>
</tr>
<tr>
<td>ALTER AGGREGATE</td>
<td>1353</td>
</tr>
<tr>
<td>ALTER COLLATION</td>
<td>1355</td>
</tr>
<tr>
<td>ALTER CONVERSION</td>
<td>1357</td>
</tr>
<tr>
<td>ALTER DATABASE</td>
<td>1359</td>
</tr>
<tr>
<td>ALTER DEFAULT PRIVILEGES</td>
<td>1362</td>
</tr>
<tr>
<td>ALTER DOMAIN</td>
<td>1365</td>
</tr>
<tr>
<td>ALTER EVENT TRIGGER</td>
<td>1368</td>
</tr>
<tr>
<td>ALTER EXTENSION</td>
<td>1369</td>
</tr>
<tr>
<td>ALTER FOREIGN DATA WRAPPER</td>
<td>1373</td>
</tr>
<tr>
<td>ALTER FOREIGN TABLE</td>
<td>1375</td>
</tr>
<tr>
<td>ALTER FUNCTION</td>
<td>1380</td>
</tr>
<tr>
<td>ALTER GROUP</td>
<td>1384</td>
</tr>
<tr>
<td>ALTER INDEX</td>
<td>1386</td>
</tr>
<tr>
<td>ALTER LANGUAGE</td>
<td>1389</td>
</tr>
<tr>
<td>ALTER LARGE OBJECT</td>
<td>1390</td>
</tr>
<tr>
<td>ALTER MATERIALIZED VIEW</td>
<td>1391</td>
</tr>
<tr>
<td>ALTER OPERATOR</td>
<td>1393</td>
</tr>
<tr>
<td>ALTER OPERATOR CLASS</td>
<td>1395</td>
</tr>
<tr>
<td>ALTER OPERATOR FAMILY</td>
<td>1396</td>
</tr>
<tr>
<td>ALTER POLICY</td>
<td>1400</td>
</tr>
<tr>
<td>ALTER PROCEDURE</td>
<td>1402</td>
</tr>
<tr>
<td>ALTER PUBLICATION</td>
<td>1405</td>
</tr>
<tr>
<td>ALTER ROLE</td>
<td>1407</td>
</tr>
<tr>
<td>ALTER ROUTINE</td>
<td>1411</td>
</tr>
<tr>
<td>ALTER RULE</td>
<td>1413</td>
</tr>
<tr>
<td>ALTER SCHEMA</td>
<td>1414</td>
</tr>
<tr>
<td>ALTER SEQUENCE</td>
<td>1415</td>
</tr>
<tr>
<td>ALTER SERVER</td>
<td>1418</td>
</tr>
<tr>
<td>ALTER STATISTICS</td>
<td>1420</td>
</tr>
<tr>
<td>ALTER SUBSCRIPTION</td>
<td>1421</td>
</tr>
<tr>
<td>ALTER SYSTEM</td>
<td>1423</td>
</tr>
<tr>
<td>ALTER TABLE</td>
<td>1425</td>
</tr>
<tr>
<td>ALTER TABLESPACE</td>
<td>1441</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH CONFIGURATION</td>
<td>1443</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH DICTIONARY</td>
<td>1445</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH PARSER</td>
<td>1447</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH TEMPLATE</td>
<td>1448</td>
</tr>
<tr>
<td>ALTER TRIGGER</td>
<td>1449</td>
</tr>
<tr>
<td>ALTER TYPE</td>
<td>1451</td>
</tr>
<tr>
<td>ALTER USER</td>
<td>1455</td>
</tr>
<tr>
<td>ALTER USER MAPPING</td>
<td>1456</td>
</tr>
<tr>
<td>ALTER VIEW</td>
<td>1457</td>
</tr>
<tr>
<td>ANALYZE</td>
<td>1459</td>
</tr>
<tr>
<td>BEGIN</td>
<td>1462</td>
</tr>
<tr>
<td>CALL</td>
<td>1464</td>
</tr>
<tr>
<td>CHECKPOINT</td>
<td>1465</td>
</tr>
<tr>
<td>CLOSE</td>
<td>1466</td>
</tr>
<tr>
<td>CLUSTER</td>
<td>1467</td>
</tr>
<tr>
<td>COMMENT</td>
<td>1469</td>
</tr>
<tr>
<td>COMMIT</td>
<td>1474</td>
</tr>
<tr>
<td>COMMIT PREPARED</td>
<td>1475</td>
</tr>
<tr>
<td>COPY</td>
<td>1476</td>
</tr>
<tr>
<td>CREATE ACCESS METHOD</td>
<td>1486</td>
</tr>
<tr>
<td>CREATE AGGREGATE</td>
<td>1487</td>
</tr>
<tr>
<td>CREATE CAST</td>
<td>1495</td>
</tr>
<tr>
<td>CREATE COLLATION</td>
<td>1499</td>
</tr>
<tr>
<td>CREATE CONVERSION</td>
<td>1501</td>
</tr>
<tr>
<td>SQL Command</td>
<td>Page</td>
</tr>
<tr>
<td>-----------------------------------</td>
<td>-------</td>
</tr>
<tr>
<td>DROP LANGUAGE</td>
<td>1665</td>
</tr>
<tr>
<td>DROP GROUP</td>
<td>1662</td>
</tr>
<tr>
<td>DROP FUNCTION</td>
<td>1660</td>
</tr>
<tr>
<td>DROP DATABASE</td>
<td>1654</td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>1653</td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>1652</td>
</tr>
<tr>
<td>DROP CAST</td>
<td>1651</td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>1650</td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>1649</td>
</tr>
<tr>
<td>DROP DATABASE</td>
<td>1648</td>
</tr>
<tr>
<td>DROP ACCESS METHOD</td>
<td>1647</td>
</tr>
<tr>
<td>CREATE SCHEMA</td>
<td>1667</td>
</tr>
<tr>
<td>CREATE SEQUENCE</td>
<td>1670</td>
</tr>
<tr>
<td>CREATE SERVER</td>
<td>1674</td>
</tr>
<tr>
<td>CREATE STATISTICS</td>
<td>1676</td>
</tr>
<tr>
<td>CREATE SUBSCRIPTION</td>
<td>1678</td>
</tr>
<tr>
<td>CREATE TABLE AS</td>
<td>1681</td>
</tr>
<tr>
<td>CREATE TABLE</td>
<td>1684</td>
</tr>
<tr>
<td>CREATE TABLE AS</td>
<td>1687</td>
</tr>
<tr>
<td>CREATE TABLESPACES</td>
<td>1690</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH DICTIONARY</td>
<td>1607</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH CONFIGURATION</td>
<td>1606</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1611</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1614</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1617</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1620</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1623</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1626</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1629</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1632</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1635</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1638</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH DICTIONARY</td>
<td>1607</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH CONFIGURATION</td>
<td>1606</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TPL</td>
<td>1611</td>
</tr>
<tr>
<td>CREATE USER MAPPING</td>
<td>1631</td>
</tr>
<tr>
<td>CREATE USER</td>
<td>1630</td>
</tr>
<tr>
<td>CREATE USER MAPPING</td>
<td>1632</td>
</tr>
<tr>
<td>CREATE VIEW</td>
<td>1633</td>
</tr>
<tr>
<td>DEALLOCATE</td>
<td>1638</td>
</tr>
<tr>
<td>DECLARE</td>
<td>1639</td>
</tr>
<tr>
<td>DELETE</td>
<td>1642</td>
</tr>
<tr>
<td>DISCARD</td>
<td>1645</td>
</tr>
<tr>
<td>DO</td>
<td>1646</td>
</tr>
<tr>
<td>DROP ACCESS METHOD</td>
<td>1648</td>
</tr>
<tr>
<td>DROP AGGREGATE</td>
<td>1649</td>
</tr>
<tr>
<td>DROP CAST</td>
<td>1651</td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>1652</td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>1653</td>
</tr>
<tr>
<td>DROP DATABASE</td>
<td>1654</td>
</tr>
<tr>
<td>DROP DOMAIN</td>
<td>1655</td>
</tr>
<tr>
<td>DROP EVENT TRIGGER</td>
<td>1656</td>
</tr>
<tr>
<td>DROP EXTENSION</td>
<td>1657</td>
</tr>
<tr>
<td>DROP FOREIGN DATA WRAPPER</td>
<td>1658</td>
</tr>
<tr>
<td>DROP FOREIGN TABLE</td>
<td>1659</td>
</tr>
<tr>
<td>DROP FUNCTION</td>
<td>1660</td>
</tr>
<tr>
<td>DROP GROUP</td>
<td>1662</td>
</tr>
<tr>
<td>DROP INDEX</td>
<td>1663</td>
</tr>
<tr>
<td>DROP LANGUAGE</td>
<td>1665</td>
</tr>
<tr>
<td>DROP MATERIALIZED VIEW</td>
<td>1666</td>
</tr>
<tr>
<td>Command</td>
<td>Page</td>
</tr>
<tr>
<td>-----------------------------</td>
<td>------</td>
</tr>
<tr>
<td>DROP OPERATOR</td>
<td>1667</td>
</tr>
<tr>
<td>DROP OPERATOR CLASS</td>
<td>1669</td>
</tr>
<tr>
<td>DROP OPERATOR FAMILY</td>
<td>1671</td>
</tr>
<tr>
<td>DROP OWNED</td>
<td>1673</td>
</tr>
<tr>
<td>DROP POLICY</td>
<td>1674</td>
</tr>
<tr>
<td>DROP PROCEDURE</td>
<td>1675</td>
</tr>
<tr>
<td>DROP PUBLICATION</td>
<td>1677</td>
</tr>
<tr>
<td>DROP ROLE</td>
<td>1678</td>
</tr>
<tr>
<td>DROP ROUTINE</td>
<td>1679</td>
</tr>
<tr>
<td>DROP RULE</td>
<td>1680</td>
</tr>
<tr>
<td>DROP SCHEMA</td>
<td>1681</td>
</tr>
<tr>
<td>DROP SEQUENCE</td>
<td>1682</td>
</tr>
<tr>
<td>DROP SERVER</td>
<td>1683</td>
</tr>
<tr>
<td>DROP STATISTICS</td>
<td>1684</td>
</tr>
<tr>
<td>DROP SUBSCRIPTION</td>
<td>1685</td>
</tr>
<tr>
<td>DROP TABLE</td>
<td>1687</td>
</tr>
<tr>
<td>DROP TABLESPACE</td>
<td>1688</td>
</tr>
<tr>
<td>DROP TEXT SEARCH CONFIGURATION</td>
<td>1689</td>
</tr>
<tr>
<td>DROP TEXT SEARCH DICTIONARY</td>
<td>1690</td>
</tr>
<tr>
<td>DROP TEXT SEARCH PARSER</td>
<td>1691</td>
</tr>
<tr>
<td>DROP TEXT SEARCH TEMPLATE</td>
<td>1692</td>
</tr>
<tr>
<td>DROP TRANSFORM</td>
<td>1693</td>
</tr>
<tr>
<td>DROP TRIGGER</td>
<td>1694</td>
</tr>
<tr>
<td>DROP TYPE</td>
<td>1695</td>
</tr>
<tr>
<td>DROP USER</td>
<td>1696</td>
</tr>
<tr>
<td>DROP USER MAPPING</td>
<td>1697</td>
</tr>
<tr>
<td>DROP VIEW</td>
<td>1698</td>
</tr>
<tr>
<td>END</td>
<td>1699</td>
</tr>
<tr>
<td>EXECUTE</td>
<td>1700</td>
</tr>
<tr>
<td>EXPLAIN</td>
<td>1701</td>
</tr>
<tr>
<td>FETCH</td>
<td>1706</td>
</tr>
<tr>
<td>GRANT</td>
<td>1710</td>
</tr>
<tr>
<td>IMPORT FOREIGN SCHEMA</td>
<td>1717</td>
</tr>
<tr>
<td>INSERT</td>
<td>1719</td>
</tr>
<tr>
<td>LISTEN</td>
<td>1726</td>
</tr>
<tr>
<td>LOAD</td>
<td>1728</td>
</tr>
<tr>
<td>LOCK</td>
<td>1729</td>
</tr>
<tr>
<td>MOVE</td>
<td>1732</td>
</tr>
<tr>
<td>NOTIFY</td>
<td>1734</td>
</tr>
<tr>
<td>PREPARE</td>
<td>1737</td>
</tr>
<tr>
<td>PREPARE TRANSACTION</td>
<td>1740</td>
</tr>
<tr>
<td>REASSIGN OWNED</td>
<td>1742</td>
</tr>
<tr>
<td>REFRESH MATERIALIZED VIEW</td>
<td>1743</td>
</tr>
<tr>
<td>REINDEX</td>
<td>1745</td>
</tr>
<tr>
<td>RELEASE SAVEPOINT</td>
<td>1748</td>
</tr>
<tr>
<td>RESET</td>
<td>1750</td>
</tr>
<tr>
<td>REVOKE</td>
<td>1751</td>
</tr>
<tr>
<td>ROLLBACK</td>
<td>1755</td>
</tr>
<tr>
<td>ROLLBACK PREPARED</td>
<td>1756</td>
</tr>
<tr>
<td>ROLLBACK TO SAVEPOINT</td>
<td>1757</td>
</tr>
<tr>
<td>SAVEPOINT</td>
<td>1759</td>
</tr>
<tr>
<td>SECURITY LABEL</td>
<td>1761</td>
</tr>
<tr>
<td>SELECT</td>
<td>1764</td>
</tr>
<tr>
<td>SELECT INTO</td>
<td>1785</td>
</tr>
<tr>
<td>SET</td>
<td>1787</td>
</tr>
<tr>
<td>SET CONSTRAINTS</td>
<td>1790</td>
</tr>
<tr>
<td>SET ROLE</td>
<td>1791</td>
</tr>
<tr>
<td>SET SESSION AUTHORIZATION</td>
<td>1793</td>
</tr>
<tr>
<td>Section</td>
<td>Title</td>
</tr>
<tr>
<td>---------</td>
<td>-------</td>
</tr>
<tr>
<td>52.9</td>
<td>pg_auth_members</td>
</tr>
<tr>
<td>52.10</td>
<td>pg_cast</td>
</tr>
<tr>
<td>52.11</td>
<td>pg_class</td>
</tr>
<tr>
<td>52.12</td>
<td>pg_collation</td>
</tr>
<tr>
<td>52.13</td>
<td>pg_constraint</td>
</tr>
<tr>
<td>52.14</td>
<td>pg_conversion</td>
</tr>
<tr>
<td>52.15</td>
<td>pg_database</td>
</tr>
<tr>
<td>52.16</td>
<td>pg_db_role_setting</td>
</tr>
<tr>
<td>52.17</td>
<td>pg_default_acl</td>
</tr>
<tr>
<td>52.18</td>
<td>pg_depend</td>
</tr>
<tr>
<td>52.19</td>
<td>pg_description</td>
</tr>
<tr>
<td>52.20</td>
<td>pg_enum</td>
</tr>
<tr>
<td>52.21</td>
<td>pg_event_trigger</td>
</tr>
<tr>
<td>52.22</td>
<td>pg_extension</td>
</tr>
<tr>
<td>52.23</td>
<td>pg_foreign_data_wrapper</td>
</tr>
<tr>
<td>52.24</td>
<td>pg_foreign_server</td>
</tr>
<tr>
<td>52.25</td>
<td>pg_foreign_table</td>
</tr>
<tr>
<td>52.26</td>
<td>pg_index</td>
</tr>
<tr>
<td>52.27</td>
<td>pg_inherits</td>
</tr>
<tr>
<td>52.28</td>
<td>pg_init_privs</td>
</tr>
<tr>
<td>52.29</td>
<td>pg_language</td>
</tr>
<tr>
<td>52.30</td>
<td>pg_largeobject</td>
</tr>
<tr>
<td>52.31</td>
<td>pg_largeobject_metadata</td>
</tr>
<tr>
<td>52.32</td>
<td>pg_namespace</td>
</tr>
<tr>
<td>52.33</td>
<td>pg_opclass</td>
</tr>
<tr>
<td>52.34</td>
<td>pg_operator</td>
</tr>
<tr>
<td>52.35</td>
<td>pg_opfamily</td>
</tr>
<tr>
<td>52.36</td>
<td>pg_partitioned_table</td>
</tr>
<tr>
<td>52.37</td>
<td>pg_pltemplate</td>
</tr>
<tr>
<td>52.38</td>
<td>pg_policy</td>
</tr>
<tr>
<td>52.39</td>
<td>pg_proc</td>
</tr>
<tr>
<td>52.40</td>
<td>pg_publication</td>
</tr>
<tr>
<td>52.41</td>
<td>pg_publication_rel</td>
</tr>
<tr>
<td>52.42</td>
<td>pg_range</td>
</tr>
<tr>
<td>52.43</td>
<td>pg_replication_origin</td>
</tr>
<tr>
<td>52.44</td>
<td>pg_rewrite</td>
</tr>
<tr>
<td>52.45</td>
<td>pg_shlabel</td>
</tr>
<tr>
<td>52.46</td>
<td>pg_sequence</td>
</tr>
<tr>
<td>52.47</td>
<td>pg_shdepend</td>
</tr>
<tr>
<td>52.48</td>
<td>pg_shdescription</td>
</tr>
<tr>
<td>52.49</td>
<td>pg_shseclabel</td>
</tr>
<tr>
<td>52.50</td>
<td>pg_statistic</td>
</tr>
<tr>
<td>52.51</td>
<td>pg_statistic_ext</td>
</tr>
<tr>
<td>52.52</td>
<td>pg_subscription</td>
</tr>
<tr>
<td>52.53</td>
<td>pg_subscription_rel</td>
</tr>
<tr>
<td>52.54</td>
<td>pg_tablespace</td>
</tr>
<tr>
<td>52.55</td>
<td>pg_transform</td>
</tr>
<tr>
<td>52.56</td>
<td>pg_trigger</td>
</tr>
<tr>
<td>52.57</td>
<td>pg_ts_config</td>
</tr>
<tr>
<td>52.58</td>
<td>pg_ts_config_map</td>
</tr>
<tr>
<td>52.59</td>
<td>pg_ts_dict</td>
</tr>
<tr>
<td>52.60</td>
<td>pg_ts_parser</td>
</tr>
<tr>
<td>52.61</td>
<td>pg_ts_template</td>
</tr>
<tr>
<td>52.62</td>
<td>pg_type</td>
</tr>
<tr>
<td>52.63</td>
<td>pg_user_mapping</td>
</tr>
<tr>
<td>52.64</td>
<td>System Views</td>
</tr>
<tr>
<td>52.65</td>
<td>pg_available_extensions</td>
</tr>
<tr>
<td>52.66</td>
<td>pg_available_extension_versions</td>
</tr>
<tr>
<td>Section</td>
<td>Title</td>
</tr>
<tr>
<td>---------</td>
<td>-------</td>
</tr>
<tr>
<td>69.5.</td>
<td>Structure of the Bootstrap BKI File</td>
</tr>
<tr>
<td>69.3.</td>
<td>BKI File Format</td>
</tr>
<tr>
<td>69.2.</td>
<td>System Catalog Initial Data</td>
</tr>
<tr>
<td>69.1.</td>
<td>System Catalog Declaration Rules</td>
</tr>
<tr>
<td>68.4.</td>
<td>Visibility Map</td>
</tr>
<tr>
<td>68.1.</td>
<td>Database File Layout</td>
</tr>
<tr>
<td>67.3.</td>
<td>Extensibility</td>
</tr>
<tr>
<td>67.2.</td>
<td>Built-in Operator Classes</td>
</tr>
<tr>
<td>67.1.</td>
<td>Introduction</td>
</tr>
<tr>
<td>66.7.</td>
<td>Examples</td>
</tr>
<tr>
<td>66.6.</td>
<td>Limitations</td>
</tr>
<tr>
<td>66.5.</td>
<td>GIN Tips and Tricks</td>
</tr>
<tr>
<td>66.4.</td>
<td>Implementation</td>
</tr>
<tr>
<td>66.3.</td>
<td>Extensibility</td>
</tr>
<tr>
<td>66.2.</td>
<td>Built-in Operator Classes</td>
</tr>
<tr>
<td>66.1.</td>
<td>Introduction</td>
</tr>
<tr>
<td>66.0</td>
<td>Index Access Method Interface Definition</td>
</tr>
</tbody>
</table>
VIII. Appendixes .............................................................................................................. 2267
A. PostgreSQL Error Codes ......................................................................................... 2292
B. Date/Time Support .................................................................................................... 2301
  B.1. Date/Time Input Interpretation ......................................................................... 2301
  B.2. Date/Time Key Words ....................................................................................... 2302
  B.3. Date/Time Configuration Files .......................................................................... 2303
  B.4. History of Units .................................................................................................. 2304
C. SQL Key Words ........................................................................................................ 2306
D. SQL Conformance ..................................................................................................... 2329
  D.1. Supported Features ............................................................................................ 2330
  D.2. Unsupported Features ......................................................................................... 2345
E. Release Notes ............................................................................................................ 2358
  E.1. Release 11.1 ........................................................................................................ 2358
  E.2. Release 11 ........................................................................................................... 2360
  E.3. Release 10.6 ......................................................................................................... 2379
  E.4. Release 10.5 ......................................................................................................... 2384
  E.5. Release 10.4 ......................................................................................................... 2389
  E.6. Release 10.3 ......................................................................................................... 2393
  E.7. Release 10.2 ......................................................................................................... 2395
  E.8. Release 10.1 ......................................................................................................... 2400
  E.9. Release 10 ........................................................................................................... 2403
  E.10. Release 9.6.11 ................................................................................................... 2423
  E.11. Release 9.6.10 ................................................................................................... 2428
  E.12. Release 9.6.9 ..................................................................................................... 2431
  E.13. Release 9.6.8 ..................................................................................................... 2434
  E.14. Release 9.6.7 ..................................................................................................... 2436
  E.15. Release 9.6.6 ..................................................................................................... 2439
  E.16. Release 9.6.5 ..................................................................................................... 2442
  E.17. Release 9.6.4 ..................................................................................................... 2444
  E.18. Release 9.6.3 ..................................................................................................... 2449
  E.19. Release 9.6.2 ..................................................................................................... 2454
  E.20. Release 9.6.1 ..................................................................................................... 2459
  E.21. Release 9.6 ......................................................................................................... 2462
  E.22. Release 9.5.15 ................................................................................................... 2480
  E.23. Release 9.5.14 ................................................................................................... 2484
  E.24. Release 9.5.13 ................................................................................................... 2487
  E.25. Release 9.5.12 ................................................................................................... 2490
  E.26. Release 9.5.11 ................................................................................................... 2492
  E.27. Release 9.5.10 ................................................................................................... 2495
  E.28. Release 9.5.9 ..................................................................................................... 2497
  E.29. Release 9.5.8 ..................................................................................................... 2499
  E.30. Release 9.5.7 ..................................................................................................... 2504
  E.31. Release 9.5.6 ..................................................................................................... 2508
  E.32. Release 9.5.5 ..................................................................................................... 2512
  E.33. Release 9.5.4 ..................................................................................................... 2516
  E.34. Release 9.5.3 ..................................................................................................... 2520
  E.35. Release 9.5.2 ..................................................................................................... 2522
  E.36. Release 9.5.1 ..................................................................................................... 2525
  E.37. Release 9.5 ......................................................................................................... 2527
  E.38. Release 9.4.20 ................................................................................................... 2540
  E.39. Release 9.4.19 ................................................................................................... 2543
  E.40. Release 9.4.18 ................................................................................................... 2546
  E.41. Release 9.4.17 ................................................................................................... 2548
  E.42. Release 9.4.16 ................................................................................................... 2550
  E.43. Release 9.4.15 ................................................................................................... 2552
E.85. Release 9.2.24 ................................................................. 2694
E.84. Release 9.2.23 ................................................................. 2696
E.83. Release 9.2.22 ................................................................. 2697
E.82. Release 9.2.21 ................................................................. 2701
E.81. Release 9.2.20 ................................................................. 2704
E.80. Release 9.2.19 ................................................................. 2707
E.79. Release 9.2.18 ................................................................. 2708
E.78. Release 9.2.17 ................................................................. 2711
E.77. Release 9.2.16 ................................................................. 2712
E.76. Release 9.2.15 ................................................................. 2713
E.75. Release 9.2.14 ................................................................. 2717
E.74. Release 9.2.13 ................................................................. 2721
E.73. Release 9.2.12 ................................................................. 2721
E.72. Release 9.2.11 ................................................................. 2722
E.71. Release 9.2.10 ................................................................. 2726
E.70. Release 9.2.9 ................................................................. 2732
E.69. Release 9.2.8 ................................................................. 2735
<table>
<thead>
<tr>
<th>Release</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>E.102. Release 9.2.7</td>
<td>2737</td>
</tr>
<tr>
<td>E.103. Release 9.2.6</td>
<td>2741</td>
</tr>
<tr>
<td>E.104. Release 9.2.5</td>
<td>2743</td>
</tr>
<tr>
<td>E.105. Release 9.2.4</td>
<td>2746</td>
</tr>
<tr>
<td>E.106. Release 9.2.3</td>
<td>2748</td>
</tr>
<tr>
<td>E.107. Release 9.2.2</td>
<td>2751</td>
</tr>
<tr>
<td>E.108. Release 9.2.1</td>
<td>2756</td>
</tr>
<tr>
<td>E.109. Release 9.2</td>
<td>2757</td>
</tr>
<tr>
<td>E.110. Release 9.1.24</td>
<td>2774</td>
</tr>
<tr>
<td>E.111. Release 9.1.23</td>
<td>2775</td>
</tr>
<tr>
<td>E.112. Release 9.1.22</td>
<td>2778</td>
</tr>
<tr>
<td>E.113. Release 9.1.21</td>
<td>2779</td>
</tr>
<tr>
<td>E.114. Release 9.1.20</td>
<td>2780</td>
</tr>
<tr>
<td>E.115. Release 9.1.19</td>
<td>2783</td>
</tr>
<tr>
<td>E.116. Release 9.1.18</td>
<td>2787</td>
</tr>
<tr>
<td>E.117. Release 9.1.17</td>
<td>2787</td>
</tr>
<tr>
<td>E.118. Release 9.1.16</td>
<td>2788</td>
</tr>
<tr>
<td>E.119. Release 9.1.15</td>
<td>2792</td>
</tr>
<tr>
<td>E.120. Release 9.1.14</td>
<td>2797</td>
</tr>
<tr>
<td>E.121. Release 9.1.13</td>
<td>2800</td>
</tr>
<tr>
<td>E.122. Release 9.1.12</td>
<td>2801</td>
</tr>
<tr>
<td>E.123. Release 9.1.11</td>
<td>2805</td>
</tr>
<tr>
<td>E.124. Release 9.1.10</td>
<td>2806</td>
</tr>
<tr>
<td>E.125. Release 9.1.9</td>
<td>2808</td>
</tr>
<tr>
<td>E.126. Release 9.1.8</td>
<td>2811</td>
</tr>
<tr>
<td>E.127. Release 9.1.7</td>
<td>2813</td>
</tr>
<tr>
<td>E.128. Release 9.1.6</td>
<td>2816</td>
</tr>
<tr>
<td>E.129. Release 9.1.5</td>
<td>2817</td>
</tr>
<tr>
<td>E.130. Release 9.1.4</td>
<td>2820</td>
</tr>
<tr>
<td>E.131. Release 9.1.3</td>
<td>2823</td>
</tr>
<tr>
<td>E.132. Release 9.1.2</td>
<td>2827</td>
</tr>
<tr>
<td>E.133. Release 9.1.1</td>
<td>2832</td>
</tr>
<tr>
<td>E.134. Release 9.1</td>
<td>2832</td>
</tr>
<tr>
<td>E.135. Release 9.0.23</td>
<td>2849</td>
</tr>
<tr>
<td>E.136. Release 9.0.22</td>
<td>2852</td>
</tr>
<tr>
<td>E.137. Release 9.0.21</td>
<td>2852</td>
</tr>
<tr>
<td>E.138. Release 9.0.20</td>
<td>2853</td>
</tr>
<tr>
<td>E.139. Release 9.0.19</td>
<td>2856</td>
</tr>
<tr>
<td>E.140. Release 9.0.18</td>
<td>2861</td>
</tr>
<tr>
<td>E.141. Release 9.0.17</td>
<td>2864</td>
</tr>
<tr>
<td>E.142. Release 9.0.16</td>
<td>2865</td>
</tr>
<tr>
<td>E.143. Release 9.0.15</td>
<td>2868</td>
</tr>
<tr>
<td>E.144. Release 9.0.14</td>
<td>2870</td>
</tr>
<tr>
<td>E.145. Release 9.0.13</td>
<td>2872</td>
</tr>
<tr>
<td>E.146. Release 9.0.12</td>
<td>2874</td>
</tr>
<tr>
<td>E.147. Release 9.0.11</td>
<td>2875</td>
</tr>
<tr>
<td>E.148. Release 9.0.10</td>
<td>2878</td>
</tr>
<tr>
<td>E.149. Release 9.0.9</td>
<td>2879</td>
</tr>
<tr>
<td>E.150. Release 9.0.8</td>
<td>2881</td>
</tr>
<tr>
<td>E.151. Release 9.0.7</td>
<td>2883</td>
</tr>
<tr>
<td>E.152. Release 9.0.6</td>
<td>2887</td>
</tr>
<tr>
<td>E.153. Release 9.0.5</td>
<td>2889</td>
</tr>
<tr>
<td>E.154. Release 9.0.4</td>
<td>2893</td>
</tr>
<tr>
<td>E.155. Release 9.0.3</td>
<td>2895</td>
</tr>
<tr>
<td>E.156. Release 9.0.2</td>
<td>2896</td>
</tr>
<tr>
<td>E.157. Release 9.0.1</td>
<td>2899</td>
</tr>
<tr>
<td>E.158. Release 9.0</td>
<td>2900</td>
</tr>
<tr>
<td>E.159. Release 8.4.22</td>
<td>2920</td>
</tr>
<tr>
<td>E.276. Release 8.0.4</td>
<td>3139</td>
</tr>
<tr>
<td>E.277. Release 8.0.3</td>
<td>3140</td>
</tr>
<tr>
<td>E.278. Release 8.0.2</td>
<td>3142</td>
</tr>
<tr>
<td>E.279. Release 8.0.1</td>
<td>3144</td>
</tr>
<tr>
<td>E.280. Release 8.0</td>
<td>3145</td>
</tr>
<tr>
<td>E.281. Release 7.4.30</td>
<td>3159</td>
</tr>
<tr>
<td>E.282. Release 7.4.29</td>
<td>3160</td>
</tr>
<tr>
<td>E.283. Release 7.4.28</td>
<td>3161</td>
</tr>
<tr>
<td>E.284. Release 7.4.27</td>
<td>3162</td>
</tr>
<tr>
<td>E.285. Release 7.4.26</td>
<td>3163</td>
</tr>
<tr>
<td>E.286. Release 7.4.25</td>
<td>3164</td>
</tr>
<tr>
<td>E.287. Release 7.4.24</td>
<td>3164</td>
</tr>
<tr>
<td>E.288. Release 7.4.23</td>
<td>3165</td>
</tr>
<tr>
<td>E.289. Release 7.4.22</td>
<td>3165</td>
</tr>
<tr>
<td>E.290. Release 7.4.21</td>
<td>3166</td>
</tr>
<tr>
<td>E.291. Release 7.4.20</td>
<td>3166</td>
</tr>
<tr>
<td>E.292. Release 7.4.19</td>
<td>3167</td>
</tr>
<tr>
<td>E.293. Release 7.4.18</td>
<td>3168</td>
</tr>
<tr>
<td>E.294. Release 7.4.17</td>
<td>3169</td>
</tr>
<tr>
<td>E.295. Release 7.4.16</td>
<td>3169</td>
</tr>
<tr>
<td>E.296. Release 7.4.15</td>
<td>3170</td>
</tr>
<tr>
<td>E.297. Release 7.4.14</td>
<td>3170</td>
</tr>
<tr>
<td>E.298. Release 7.4.13</td>
<td>3171</td>
</tr>
<tr>
<td>E.299. Release 7.4.12</td>
<td>3172</td>
</tr>
<tr>
<td>E.300. Release 7.4.11</td>
<td>3173</td>
</tr>
<tr>
<td>E.301. Release 7.4.10</td>
<td>3173</td>
</tr>
<tr>
<td>E.302. Release 7.4.9</td>
<td>3174</td>
</tr>
<tr>
<td>E.303. Release 7.4.8</td>
<td>3175</td>
</tr>
<tr>
<td>E.304. Release 7.4.7</td>
<td>3177</td>
</tr>
<tr>
<td>E.305. Release 7.4.6</td>
<td>3178</td>
</tr>
<tr>
<td>E.306. Release 7.4.5</td>
<td>3179</td>
</tr>
<tr>
<td>E.307. Release 7.4.4</td>
<td>3179</td>
</tr>
<tr>
<td>E.308. Release 7.4.3</td>
<td>3180</td>
</tr>
<tr>
<td>E.309. Release 7.4.2</td>
<td>3181</td>
</tr>
<tr>
<td>E.310. Release 7.4.1</td>
<td>3182</td>
</tr>
<tr>
<td>E.311. Release 7.4</td>
<td>3184</td>
</tr>
<tr>
<td>E.312. Release 7.3.21</td>
<td>3199</td>
</tr>
<tr>
<td>E.313. Release 7.3.20</td>
<td>3200</td>
</tr>
<tr>
<td>E.314. Release 7.3.19</td>
<td>3200</td>
</tr>
<tr>
<td>E.315. Release 7.3.18</td>
<td>3201</td>
</tr>
<tr>
<td>E.316. Release 7.3.17</td>
<td>3201</td>
</tr>
<tr>
<td>E.317. Release 7.3.16</td>
<td>3201</td>
</tr>
<tr>
<td>E.318. Release 7.3.15</td>
<td>3202</td>
</tr>
<tr>
<td>E.319. Release 7.3.14</td>
<td>3203</td>
</tr>
<tr>
<td>E.320. Release 7.3.13</td>
<td>3203</td>
</tr>
<tr>
<td>E.321. Release 7.3.12</td>
<td>3204</td>
</tr>
<tr>
<td>E.322. Release 7.3.11</td>
<td>3204</td>
</tr>
<tr>
<td>E.323. Release 7.3.10</td>
<td>3205</td>
</tr>
<tr>
<td>E.324. Release 7.3.9</td>
<td>3206</td>
</tr>
<tr>
<td>E.325. Release 7.3.8</td>
<td>3207</td>
</tr>
<tr>
<td>E.326. Release 7.3.7</td>
<td>3208</td>
</tr>
<tr>
<td>E.327. Release 7.3.6</td>
<td>3208</td>
</tr>
<tr>
<td>E.328. Release 7.3.5</td>
<td>3209</td>
</tr>
<tr>
<td>E.329. Release 7.3.4</td>
<td>3210</td>
</tr>
<tr>
<td>E.330. Release 7.3.3</td>
<td>3210</td>
</tr>
<tr>
<td>E.331. Release 7.3.2</td>
<td>3213</td>
</tr>
<tr>
<td>E.332. Release 7.3.1</td>
<td>3214</td>
</tr>
<tr>
<td>E.333. Release 7.3</td>
<td>3214</td>
</tr>
</tbody>
</table>
List of Figures

9.1. XSLT Stylesheet for Converting SQL/XML Output to HTML ............................................. 284
60.1. Structured Diagram of a Genetic Algorithm ................................................................. 2185
List of Tables

4.1. Backslash Escape Sequences .......................................................... 35
4.2. Operator Precedence (highest to lowest) ........................................... 40
8.1. Data Types .................................................................................. 131
8.2. Numeric Types ............................................................................. 132
8.3. Monetary Types ............................................................................ 137
8.4. Character Types ............................................................................ 138
8.5. Special Character Types ................................................................. 140
8.6. Binary Data Types ......................................................................... 140
8.7. bytea Literal Escaped Octets ............................................................. 141
8.8. bytea Output Escaped Octets ......................................................... 142
8.9. Date/Time Types ........................................................................... 142
8.10. Date Input ................................................................................... 144
8.11. Time Input .................................................................................. 144
8.12. Time Zone Input ........................................................................... 145
8.13. Special Date/Time Inputs ............................................................... 146
8.14. Date/Time Output Styles ............................................................... 147
8.15. Date Order Conventions ............................................................... 147
8.16. ISO 8601 Interval Unit Abbreviations ............................................ 149
8.17. Interval Input ................................................................................ 150
8.18. Interval Output Style Examples ...................................................... 151
8.19. Boolean Data Type ....................................................................... 151
8.20. Geometric Types .......................................................................... 154
8.21. Network Address Types ............................................................... 157
8.22. cidr Type Input Examples ............................................................. 157
8.23. JSON primitive types and corresponding PostgreSQL types .......... 166
8.24. Object Identifier Types ................................................................. 195
8.25. Pseudo-Types .............................................................................. 196
9.1. Comparison Operators ................................................................. 198
9.2. Comparison Predicates .................................................................. 199
9.3. Comparison Functions .................................................................. 201
9.4. Mathematical Operators ............................................................... 201
9.5. Mathematical Functions ............................................................... 202
9.6. Random Functions ....................................................................... 204
9.7. Trigonometric Functions ............................................................... 204
9.8. SQL String Functions and Operators .............................................. 205
9.9. Other String Functions ................................................................. 206
9.10. Built-in Conversions .................................................................... 213
9.11. SQL Binary String Functions and Operators .................................. 219
9.12. Other Binary String Functions ...................................................... 220
9.13. Bit String Operators .................................................................... 221
9.14. Regular Expression Match Operators ............................................ 225
9.15. Regular Expression Atoms ........................................................... 229
9.16. Regular Expression Quantifiers ...................................................... 229
9.17. Regular Expression Constraints .................................................... 230
9.18. Regular Expression Character-entry Escapes .................................. 232
9.19. Regular Expression Class-shorthand Escapes ................................ 233
9.20. Regular Expression Constraint Escapes .......................................... 233
9.21. Regular Expression Back References ............................................ 233
9.22. ARE Embedded-option Letters ..................................................... 234
9.23. Formatting Functions ................................................................... 237
9.24. Template Patterns for Date/Time Formatting .................................. 238
9.25. Template Pattern Modifiers for Date/Time Formatting .................... 240
9.26. Template Patterns for Numeric Formatting .................................... 242
9.27. Template Pattern Modifiers for Numeric Formatting ...................... 243
9.28. to_char Examples ........................................................................ 243
<table>
<thead>
<tr>
<th>Section Number</th>
<th>Section Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>9.87</td>
<td>Index Maintenance Functions</td>
<td>351</td>
</tr>
<tr>
<td>9.88</td>
<td>Generic File Access Functions</td>
<td>351</td>
</tr>
<tr>
<td>9.89</td>
<td>Advisory Lock Functions</td>
<td>353</td>
</tr>
<tr>
<td>9.90</td>
<td>Table Rewrite information</td>
<td>358</td>
</tr>
<tr>
<td>12.1</td>
<td>Default Parser's Token Types</td>
<td>404</td>
</tr>
<tr>
<td>13.1</td>
<td>Transaction Isolation Levels</td>
<td>427</td>
</tr>
<tr>
<td>13.2</td>
<td>Conflicting Lock Modes</td>
<td>434</td>
</tr>
<tr>
<td>13.3</td>
<td>Conflicting Row-level Locks</td>
<td>435</td>
</tr>
<tr>
<td>18.1</td>
<td>System V IPC Parameters</td>
<td>510</td>
</tr>
<tr>
<td>18.2</td>
<td>SSL Server File Usage</td>
<td>525</td>
</tr>
<tr>
<td>19.1</td>
<td>Message Severity Levels</td>
<td>566</td>
</tr>
<tr>
<td>19.2</td>
<td>Short Option Key</td>
<td>592</td>
</tr>
<tr>
<td>21.1</td>
<td>Default Roles</td>
<td>618</td>
</tr>
<tr>
<td>23.1</td>
<td>PostgreSQL Character Sets</td>
<td>635</td>
</tr>
<tr>
<td>23.2</td>
<td>Client/Server Character Set Conversions</td>
<td>639</td>
</tr>
<tr>
<td>28.2</td>
<td>VACUUM phases</td>
<td>727</td>
</tr>
<tr>
<td>28.17</td>
<td>Per-Backend Statistics Functions</td>
<td>725</td>
</tr>
<tr>
<td>28.16</td>
<td>Collected Statistics Views</td>
<td>697</td>
</tr>
<tr>
<td>28.15</td>
<td>pg_stat_activity View</td>
<td>698</td>
</tr>
<tr>
<td>28.14</td>
<td>wait_event Description</td>
<td>702</td>
</tr>
<tr>
<td>28.13</td>
<td>pg_stat_replication View</td>
<td>712</td>
</tr>
<tr>
<td>28.12</td>
<td>pg_stat_database View</td>
<td>718</td>
</tr>
<tr>
<td>28.11</td>
<td>pg_stat_database_conflicts View</td>
<td>720</td>
</tr>
<tr>
<td>28.10</td>
<td>pg_stat_all_tables View</td>
<td>720</td>
</tr>
<tr>
<td>28.9</td>
<td>pg_stat_archiver View</td>
<td>717</td>
</tr>
<tr>
<td>28.8</td>
<td>pg_stat_ssl View</td>
<td>717</td>
</tr>
<tr>
<td>28.7</td>
<td>pg_stat_subscription View</td>
<td>716</td>
</tr>
<tr>
<td>28.6</td>
<td>pg_stat_wal_receiver View</td>
<td>715</td>
</tr>
<tr>
<td>28.5</td>
<td>pg_stat_replication View</td>
<td>712</td>
</tr>
<tr>
<td>28.4</td>
<td>wait_event Description</td>
<td>702</td>
</tr>
<tr>
<td>28.3</td>
<td>pg_stat_activity View</td>
<td>698</td>
</tr>
<tr>
<td>28.2</td>
<td>Collected Statistics Views</td>
<td>697</td>
</tr>
<tr>
<td>28.1</td>
<td>Dynamic Statistics Views</td>
<td>696</td>
</tr>
<tr>
<td>23.2</td>
<td>Client/Server Character Set Conversions</td>
<td>639</td>
</tr>
<tr>
<td>23.1</td>
<td>PostgreSQL Character Sets</td>
<td>635</td>
</tr>
<tr>
<td>23.1</td>
<td>PostgreSQL Character Sets</td>
<td>635</td>
</tr>
<tr>
<td>21.1</td>
<td>Default Roles</td>
<td>618</td>
</tr>
<tr>
<td>18.2</td>
<td>SSL Server File Usage</td>
<td>525</td>
</tr>
<tr>
<td>13.2</td>
<td>Conflicting Lock Modes</td>
<td>434</td>
</tr>
<tr>
<td>13.1</td>
<td>Transaction Isolation Levels</td>
<td>427</td>
</tr>
<tr>
<td>12.1</td>
<td>Default Parser's Token Types</td>
<td>404</td>
</tr>
<tr>
<td>9.88</td>
<td>Generic File Access Functions</td>
<td>351</td>
</tr>
<tr>
<td>9.89</td>
<td>Advisory Lock Functions</td>
<td>353</td>
</tr>
<tr>
<td>9.90</td>
<td>Table Rewrite information</td>
<td>358</td>
</tr>
</tbody>
</table>

xxx
52.45. pg_seclabel Columns ......................................................... 2046
52.46. pg_sequence Columns ..................................................... 2047
52.47. pg_shdepend Columns ..................................................... 2047
52.48. pg_shdescription Columns .............................................. 2048
52.49. pg_shseclabel Columns .................................................. 2049
52.50. pg_statistic Columns .................................................... 2050
52.51. pg_statistic_ext Columns ............................................... 2051
52.52. pg_subscription Columns .............................................. 2052
52.53. pg_subscription_rel Columns ........................................... 2053
52.54. pg_tablespace Columns ................................................ 2053
52.55. pg_transform Columns .................................................. 2054
52.56. pg_trigger Columns ..................................................... 2054
52.57. pg_ts_config Columns .................................................. 2056
52.58. pg_ts_config_map Columns ............................................ 2057
52.59. pg_ts_ddict Columns .................................................... 2057
52.60. pg_ts_parser Columns .................................................. 2058
52.61. pg_ts_template Columns .............................................. 2058
52.62. pg_type Columns ......................................................... 2059
52.63. typcategory Codes ...................................................... 2064
52.64. pg_user_mapping Columns ............................................. 2065
52.65. System Views ............................................................ 2066
52.66. pg_available_extensions Columns .................................... 2066
52.67. pg_available_extension_versions Columns ......................... 2067
52.68. pg_config Columns ...................................................... 2067
52.69. pg_containers Columns ................................................ 2068
52.70. pg_file_settings Columns ........................................... 2069
52.71. pg_group Columns ....................................................... 2069
52.72. pg_hba_file_rules Columns ........................................... 2070
52.73. pg_indexes Columns .................................................... 2070
52.74. pg_locks Columns ....................................................... 2071
52.75. pg_matviews Columns ................................................... 2074
52.76. pg_policies Columns .................................................... 2074
52.77. pg_prepared_statements Columns ................................... 2075
52.78. pg_prepared_xacts Columns ........................................... 2076
52.79. pg_publication_tables Columns ..................................... 2076
52.80. pg_replication_origin_status Columns .............................. 2076
52.81. pg_replication_slots Columns ....................................... 2077
52.82. pg_roles Columns ....................................................... 2078
52.83. pg_rules Columns ....................................................... 2079
52.84. pg_seclabels Columns .................................................. 2080
52.85. pg_sequences Columns ................................................ 2080
52.86. pg_settings Columns .................................................. 2081
52.87. pg_shadow Columns ..................................................... 2083
52.88. pg_stats Columns ....................................................... 2084
52.89. pg_tables Columns ..................................................... 2087
52.90. pg_timezone_abbrs Columns .......................................... 2087
52.91. pg_timezone_names Columns ......................................... 2087
52.92. pg_user Columns ......................................................... 2088
52.93. pg_user_mappings Columns .......................................... 2088
52.94. pg_views Columns ....................................................... 2089
64.1. Built-in GiST Operator Classes ..................................... 2207
65.1. Built-in SP-GiST Operator Classes .................................. 2219
66.1. Built-in GIN Operator Classes ...................................... 2230
67.1. Built-in BRIN Operator Classes ..................................... 2237
67.2. Function and Support Numbers for Minmax Operator Classes .......... 2239
67.3. Function and Support Numbers for Inclusion Operator Classes .......... 2239
68.1. Contents of PGDATA ...................................................... 2241
68.2. Page Layout ................................................................ 2247
List of Examples

8.1. Using the Character Types ................................................................. 139
8.2. Using the boolean Type ................................................................. 152
8.3. Using the Bit String Types .......................................................... 159
10.1. Factorial Operator Type Resolution ............................................. 361
10.2. String Concatenation Operator Type Resolution .......................... 362
10.3. Absolute-Value and Negation Operator Type Resolution ............... 363
10.4. Array Inclusion Operator Type Resolution .................................... 363
10.5. Custom Operator on a Domain Type ............................................. 363
10.6. Rounding Function Argument Type Resolution ........................... 366
10.7. Variadic Function Resolution ...................................................... 366
10.8. Substring Function Type Resolution ............................................. 367
10.9. character Storage Type Conversion ........................................... 368
10.10. Type Resolution with Underspecified Types in a Union ........... 369
10.11. Type Resolution in a Simple Union ............................................ 369
10.12. Type Resolution in a Transposed Union .................................... 369
10.13. Type Resolution in a Nested Union ........................................... 370
11.1. Setting up a Partial Index to Exclude Common Values ............... 378
11.2. Setting up a Partial Index to Exclude Uninteresting Values .......... 378
11.3. Setting up a Partial Unique Index .............................................. 379
20.1. Example pg_hba.conf Entries ...................................................... 599
20.2. An Example pg_ident.conf File .................................................... 602
34.1. libpq Example Program 1 .......................................................... 841
34.2. libpq Example Program 2 .......................................................... 844
34.3. libpq Example Program 3 .......................................................... 844
35.1. Large Objects with libpq Example Program ................................ 856
36.1. Example SQLDA Program .......................................................... 908
36.2. ECPG Program Accessing Large Objects ................................... 922
42.1. Manual Installation of PL/Perl .................................................... 1157
43.1. Quoting Values In Dynamic Queries .......................................... 1172
43.2. Exceptions with UPDATE/INSERT ............................................. 1187
43.3. A PL/pgSQL Trigger Function .................................................... 1201
43.4. A PL/pgSQL Trigger Function For Auditing ............................... 1202
43.5. A PL/pgSQL View Trigger Function For Auditing ....................... 1203
43.6. A PL/pgSQL Trigger Function For Maintaining A Summary Table .... 1204
43.7. Auditing with Transition Tables ................................................ 1206
43.8. A PL/pgSQL Event Trigger Function .......................................... 1208
43.9. Porting a Simple Function from PL/SQL to PL/pgSQL .................. 1215
43.10. Porting a Function that Creates Another Function from PL/SQL to PL/pgSQL .................................................. 1216
43.11. Porting a Procedure With String Manipulation and OUT Parameters from PL/SQL to PL/pgSQL .................................................. 1217
43.12. Porting a Procedure from PL/SQL to PL/pgSQL ......................... 1219
F.1. Create a Foreign Table for PostgreSQL CSV Logs ....................... 3349
Preface

This book is the official documentation of PostgreSQL. It has been written by the PostgreSQL developers and other volunteers in parallel to the development of the PostgreSQL software. It describes all the functionality that the current version of PostgreSQL officially supports.

To make the large amount of information about PostgreSQL manageable, this book has been organized in several parts. Each part is targeted at a different class of users, or at users in different stages of their PostgreSQL experience:

• Part I is an informal introduction for new users.
• Part II documents the SQL query language environment, including data types and functions, as well as user-level performance tuning. Every PostgreSQL user should read this.
• Part III describes the installation and administration of the server. Everyone who runs a PostgreSQL server, be it for private use or for others, should read this part.
• Part IV describes the programming interfaces for PostgreSQL client programs.
• Part V contains information for advanced users about the extensibility capabilities of the server. Topics include user-defined data types and functions.
• Part VI contains reference information about SQL commands, client and server programs. This part supports the other parts with structured information sorted by command or program.
• Part VII contains assorted information that might be of use to PostgreSQL developers.

1. What is PostgreSQL?

PostgreSQL is an object-relational database management system (ORDBMS) based on POSTGRES, Version 4.2, developed at the University of California at Berkeley Computer Science Department. POSTGRES pioneered many concepts that only became available in some commercial database systems much later.

PostgreSQL is an open-source descendant of this original Berkeley code. It supports a large part of the SQL standard and offers many modern features:

• complex queries
• foreign keys
• triggers
• updatable views
• transactional integrity
• multiversion concurrency control

Also, PostgreSQL can be extended by the user in many ways, for example by adding new

• data types
• functions
• operators
• aggregate functions
• index methods
• procedural languages

And because of the liberal license, PostgreSQL can be used, modified, and distributed by anyone free of charge for any purpose, be it private, commercial, or academic.

2. A Brief History of PostgreSQL

1 http://db.cs.berkeley.edu/postres.html
The object-relational database management system now known as PostgreSQL is derived from the POSTGRES package written at the University of California at Berkeley. With over two decades of development behind it, PostgreSQL is now the most advanced open-source database available anywhere.

2.1. The Berkeley POSTGRES Project

The POSTGRES project, led by Professor Michael Stonebraker, was sponsored by the Defense Advanced Research Projects Agency (DARPA), the Army Research Office (ARO), the National Science Foundation (NSF), and ESL, Inc. The implementation of POSTGRES began in 1986. The initial concepts for the system were presented in [ston86], and the definition of the initial data model appeared in [rowe87]. The design of the rule system at that time was described in [ston87a]. The rationale and architecture of the storage manager were detailed in [ston87b].

POSTGRES has undergone several major releases since then. The first “demoware” system became operational in 1987 and was shown at the 1988 ACM-SIGMOD Conference. Version 1, described in [ston90a], was released to a few external users in June 1989. In response to a critique of the first rule system ([ston89]), the rule system was redesigned ([ston90b]), and Version 2 was released in June 1990 with the new rule system. Version 3 appeared in 1991 and added support for multiple storage managers, an improved query executor, and a rewritten rule system. For the most part, subsequent releases until Postgres95 (see below) focused on portability and reliability.

POSTGRES has been used to implement many different research and production applications. These include: a financial data analysis system, a jet engine performance monitoring package, an asteroid tracking database, a medical information database, and several geographic information systems. POSTGRES has also been used as an educational tool at several universities. Finally, Illustra Information Technologies (later merged into Informix) picked up the code and commercialized it. In late 1992, POSTGRES became the primary data manager for the Sequoia 2000 scientific computing project.

The size of the external user community nearly doubled during 1993. It became increasingly obvious that maintenance of the prototype code and support was taking up large amounts of time that should have been devoted to database research. In an effort to reduce this support burden, the Berkeley POSTGRES project officially ended with Version 4.2.

2.2. Postgres95

In 1994, Andrew Yu and Jolly Chen added an SQL language interpreter to POSTGRES. Under a new name, Postgres95 was subsequently released to the web to find its own way in the world as an open-source descendant of the original POSTGRES Berkeley code.

Postgres95 code was completely ANSI C and trimmed in size by 25%. Many internal changes improved performance and maintainability. Postgres95 release 1.0.x ran about 30-50% faster on the Wisconsin Benchmark compared to POSTGRES, Version 4.2. Apart from bug fixes, the following were the major enhancements:

- The query language PostQUEL was replaced with SQL (implemented in the server). (Interface library libpq was named after PostQUEL.) Subqueries were not supported until PostgreSQL (see below), but they could be imitated in Postgres95 with user-defined SQL functions. Aggregate functions were re-implemented. Support for the GROUP BY query clause was also added.

- A new program (psql) was provided for interactive SQL queries, which used GNU Readline. This largely superseded the old monitor program.

- A new front-end library, libpgtcl, supported Tcl-based clients. A sample shell, pgtclsh, provided new Tcl commands to interface Tcl programs with the Postgres95 server.

---

2 https://www.ibm.com/analytics/informix
3 https://www.ibm.com/
4 http://meteora.ucsd.edu/s2k/s2k_home.html
• The large-object interface was overhauled. The inversion large objects were the only mechanism for storing large objects. (The inversion file system was removed.)

• The instance-level rule system was removed. Rules were still available as rewrite rules.

• A short tutorial introducing regular SQL features as well as those of Postgres95 was distributed with the source code

• GNU make (instead of BSD make) was used for the build. Also, Postgres95 could be compiled with an unpatched GCC (data alignment of doubles was fixed).

2.3. PostgreSQL

By 1996, it became clear that the name “Postgres95” would not stand the test of time. We chose a new name, PostgreSQL, to reflect the relationship between the original POSTGRES and the more recent versions with SQL capability. At the same time, we set the version numbering to start at 6.0, putting the numbers back into the sequence originally begun by the Berkeley POSTGRES project.

Many people continue to refer to PostgreSQL as “Postgres” (now rarely in all capital letters) because of tradition or because it is easier to pronounce. This usage is widely accepted as a nickname or alias.

The emphasis during development of Postgres95 was on identifying and understanding existing problems in the server code. With PostgreSQL, the emphasis has shifted to augmenting features and capabilities, although work continues in all areas.

Details about what has happened in PostgreSQL since then can be found in Appendix E.

3. Conventions

The following conventions are used in the synopsis of a command: brackets ([ and ]) indicate optional parts. (In the synopsis of a Tcl command, question marks (?) are used instead, as is usual in Tcl.) Braces ({ and }) and vertical lines (|) indicate that you must choose one alternative. Dots (... ) mean that the preceding element can be repeated.

Where it enhances the clarity, SQL commands are preceded by the prompt =>, and shell commands are preceded by the prompt $. Normally, prompts are not shown, though.

An administrator is generally a person who is in charge of installing and running the server. A user could be anyone who is using, or wants to use, any part of the PostgreSQL system. These terms should not be interpreted too narrowly; this book does not have fixed presumptions about system administration procedures.

4. Further Information

Besides the documentation, that is, this book, there are other resources about PostgreSQL:

Wiki

The PostgreSQL wiki\(^5\) contains the project’s FAQ\(^6\) (Frequently Asked Questions) list, TODO\(^7\) list, and detailed information about many more topics.

Web Site

The PostgreSQL web site\(^8\) carries details on the latest release and other information to make your work or play with PostgreSQL more productive.

---

\(^5\) https://wiki.postgresql.org
\(^6\) https://wiki.postgresql.org/wiki/Frequently_Asked_Questions
\(^7\) https://wiki.postgresql.org/wiki/TODO
\(^8\) https://www.postgresql.org
Mailing Lists

The mailing lists are a good place to have your questions answered, to share experiences with other users, and to contact the developers. Consult the PostgreSQL web site for details.

Yourself!

PostgreSQL is an open-source project. As such, it depends on the user community for ongoing support. As you begin to use PostgreSQL, you will rely on others for help, either through the documentation or through the mailing lists. Consider contributing your knowledge back. Read the mailing lists and answer questions. If you learn something which is not in the documentation, write it up and contribute it. If you add features to the code, contribute them.

5. Bug Reporting Guidelines

When you find a bug in PostgreSQL we want to hear about it. Your bug reports play an important part in making PostgreSQL more reliable because even the utmost care cannot guarantee that every part of PostgreSQL will work on every platform under every circumstance.

The following suggestions are intended to assist you in forming bug reports that can be handled in an effective fashion. No one is required to follow them but doing so tends to be to everyone's advantage.

We cannot promise to fix every bug right away. If the bug is obvious, critical, or affects a lot of users, chances are good that someone will look into it. It could also happen that we tell you to update to a newer version to see if the bug happens there. Or we might decide that the bug cannot be fixed before some major rewrite we might be planning is done. Or perhaps it is simply too hard and there are more important things on the agenda. If you need help immediately, consider obtaining a commercial support contract.

5.1. Identifying Bugs

Before you report a bug, please read and re-read the documentation to verify that you can really do whatever it is you are trying. If it is not clear from the documentation whether you can do something or not, please report that too; it is a bug in the documentation. If it turns out that a program does something different from what the documentation says, that is a bug. That might include, but is not limited to, the following circumstances:

- A program terminates with a fatal signal or an operating system error message that would point to a problem in the program. (A counterexample might be a “disk full” message, since you have to fix that yourself.)
- A program produces the wrong output for any given input.
- A program refuses to accept valid input (as defined in the documentation).
- A program accepts invalid input without a notice or error message. But keep in mind that your idea of invalid input might be our idea of an extension or compatibility with traditional practice.
- PostgreSQL fails to compile, build, or install according to the instructions on supported platforms.

Here “program” refers to any executable, not only the backend process.

Being slow or resource-hogging is not necessarily a bug. Read the documentation or ask on one of the mailing lists for help in tuning your applications. Failing to comply to the SQL standard is not necessarily a bug either, unless compliance for the specific feature is explicitly claimed.

Before you continue, check on the TODO list and in the FAQ to see if your bug is already known. If you cannot decode the information on the TODO list, report your problem. The least we can do is make the TODO list clearer.
5.2. What to Report

The most important thing to remember about bug reporting is to state all the facts and only facts. Do not speculate what you think went wrong, what “it seemed to do”, or which part of the program has a fault. If you are not familiar with the implementation you would probably guess wrong and not help us a bit. And even if you are, educated explanations are a great supplement to but no substitute for facts. If we are going to fix the bug we still have to see it happen for ourselves first. Reporting the bare facts is relatively straightforward (you can probably copy and paste them from the screen) but all too often important details are left out because someone thought it does not matter or the report would be understood anyway.

The following items should be contained in every bug report:

- The exact sequence of steps from program start-up necessary to reproduce the problem. This should be self-contained; it is not enough to send in a bare SELECT statement without the preceding CREATE TABLE and INSERT statements, if the output should depend on the data in the tables. We do not have the time to reverse-engineer your database schema, and if we are supposed to make up our own data we would probably miss the problem.

The best format for a test case for SQL-related problems is a file that can be run through the psql frontend that shows the problem. (Be sure to not have anything in your ~/.psqlrc start-up file.) An easy way to create this file is to use pg_dump to dump out the table declarations and data needed to set the scene, then add the problem query. You are encouraged to minimize the size of your example, but this is not absolutely necessary. If the bug is reproducible, we will find it either way.

If your application uses some other client interface, such as PHP, then please try to isolate the offending queries. We will probably not set up a web server to reproduce your problem. In any case remember to provide the exact input files; do not guess that the problem happens for “large files” or “midsize databases”, etc. since this information is too inexact to be of use.

- The output you got. Please do not say that it “didn’t work” or “crashed”. If there is an error message, show it, even if you do not understand it. If the program terminates with an operating system error, say which. If nothing at all happens, say so. Even if the result of your test case is a program crash or otherwise obvious it might not happen on our platform. The easiest thing is to copy the output from the terminal, if possible.

Note

If you are reporting an error message, please obtain the most verbose form of the message. In psql, say \set VERBOSITY verbose beforehand. If you are extracting the message from the server log, set the run-time parameter log_error_verbosity to verbose so that all details are logged.

Note

In case of fatal errors, the error message reported by the client might not contain all the information available. Please also look at the log output of the database server. If you do not keep your server’s log output, this would be a good time to start doing so.

- The output you expected is very important to state. If you just write “This command gives me that output.” or “This is not what I expected.”, we might run it ourselves, scan the output, and think it looks OK and is exactly what we expected. We should not have to spend the time to decode the exact semantics behind your commands. Especially refrain from merely saying that “This is not what SQL says/Oracle does.” Digging out the correct behavior from SQL is not a fun undertaking,
nor do we all know how all the other relational databases out there behave. (If your problem is a program crash, you can obviously omit this item.)

- Any command line options and other start-up options, including any relevant environment variables or configuration files that you changed from the default. Again, please provide exact information. If you are using a prepackaged distribution that starts the database server at boot time, you should try to find out how that is done.

- Anything you did at all differently from the installation instructions.

- The PostgreSQL version. You can run the command `SELECT version();` to find out the version of the server you are connected to. Most executable programs also support a `--version` option; at least `postgres --version` and `psql --version` should work. If the function or the options do not exist then your version is more than old enough to warrant an upgrade. If you run a prepackaged version, such as RPMs, say so, including any subversion the package might have. If you are talking about a Git snapshot, mention that, including the commit hash.

If your version is older than 11.1 we will almost certainly tell you to upgrade. There are many bug fixes and improvements in each new release, so it is quite possible that a bug you have encountered in an older release of PostgreSQL has already been fixed. We can only provide limited support for sites using older releases of PostgreSQL: if you require more than we can provide, consider acquiring a commercial support contract.

- Platform information. This includes the kernel name and version, C library, processor, memory information, and so on. In most cases it is sufficient to report the vendor and version, but do not assume everyone knows what exactly “Debian” contains or that everyone runs on i386s. If you have installation problems then information about the toolchain on your machine (compiler, make, and so on) is also necessary.

Do not be afraid if your bug report becomes rather lengthy. That is a fact of life. It is better to report everything the first time than us having to squeeze the facts out of you. On the other hand, if your input files are huge, it is fair to ask first whether somebody is interested in looking into it. Here is an article\(^9\) that outlines some more tips on reporting bugs.

Do not spend all your time to figure out which changes in the input make the problem go away. This will probably not help solving it. If it turns out that the bug cannot be fixed right away, you will still have time to find and share your work-around. Also, once again, do not waste your time guessing why the bug exists. We will find that out soon enough.

When writing a bug report, please avoid confusing terminology. The software package in total is called “PostgreSQL”, sometimes “Postgres” for short. If you are specifically talking about the backend process, mention that, do not just say “PostgreSQL crashes”. A crash of a single backend process is quite different from crash of the parent “postgres” process; please don’t say “the server crashed” when you mean a single backend process went down, nor vice versa. Also, client programs such as the interactive frontend “psql” are completely separate from the backend. Please try to be specific about whether the problem is on the client or server side.

### 5.3. Where to Report Bugs

In general, send bug reports to the bug report mailing list at `<pgsql-bugs@postgresql.org>`.

You are requested to use a descriptive subject for your email message, perhaps parts of the error message.

Another method is to fill in the bug report web-form available at the project’s web site\(^{10}\). Entering a bug report this way causes it to be mailed to the `<pgsql-bugs@postgresql.org>` mailing list.

---

\(^9\) [https://www.chiark.greenend.org.uk/~sgtatham/bugs.html](https://www.chiark.greenend.org.uk/~sgtatham/bugs.html)

\(^{10}\) [https://www.postgresql.org/](https://www.postgresql.org/)
If your bug report has security implications and you'd prefer that it not become immediately visible in public archives, don't send it to `pgsql-bugs`. Security issues can be reported privately to `<security@postgresql.org>`.

Do not send bug reports to any of the user mailing lists, such as `<pgsql-sql@postgresql.org>` or `<pgsql-general@postgresql.org>`. These mailing lists are for answering user questions, and their subscribers normally do not wish to receive bug reports. More importantly, they are unlikely to fix them.

Also, please do not send reports to the developers' mailing list `<pgsql-hackers@postgresql.org>`. This list is for discussing the development of PostgreSQL, and it would be nice if we could keep the bug reports separate. We might choose to take up a discussion about your bug report on `<pgsql-hackers>`, if the problem needs more review.

If you have a problem with the documentation, the best place to report it is the documentation mailing list `<pgsql-docs@postgresql.org>`. Please be specific about what part of the documentation you are unhappy with.

If your bug is a portability problem on a non-supported platform, send mail to `<pgsql-hackers@postgresql.org>`, so we (and you) can work on porting PostgreSQL to your platform.

---

**Note**

Due to the unfortunate amount of spam going around, all of the above email addresses are closed mailing lists. That is, you need to be subscribed to a list to be allowed to post on it. (You need not be subscribed to use the bug-report web form, however.) If you would like to send mail but do not want to receive list traffic, you can subscribe and set your subscription option to `nomail`. For more information send mail to `<majordomo@postgresql.org>` with the single word `help` in the body of the message.
Part I. Tutorial

Welcome to the PostgreSQL Tutorial. The following few chapters are intended to give a simple introduction to PostgreSQL, relational database concepts, and the SQL language to those who are new to any one of these aspects. We only assume some general knowledge about how to use computers. No particular Unix or programming experience is required. This part is mainly intended to give you some hands-on experience with important aspects of the PostgreSQL system. It makes no attempt to be a complete or thorough treatment of the topics it covers.

After you have worked through this tutorial you might want to move on to reading Part II to gain a more formal knowledge of the SQL language, or Part IV for information about developing applications for PostgreSQL. Those who set up and manage their own server should also read Part III.
Table of Contents

1. Getting Started .................................................................................................................. 3
   1.1. Installation .................................................................................................................... 3
   1.2. Architectural Fundamentals ....................................................................................... 3
   1.3. Creating a Database .................................................................................................... 3
   1.4. Accessing a Database ................................................................................................. 5
2. The SQL Language ........................................................................................................... 7
   2.1. Introduction .................................................................................................................. 7
   2.2. Concepts ...................................................................................................................... 7
   2.3. Creating a New Table .................................................................................................. 7
   2.4. Populating a Table With Rows .................................................................................... 8
   2.5. Querying a Table ......................................................................................................... 9
   2.6. Joins Between Tables ................................................................................................ 11
   2.7. Aggregate Functions ................................................................................................. 13
   2.8. Updates ....................................................................................................................... 15
   2.9. Deletions ..................................................................................................................... 15
3. Advanced Features ........................................................................................................... 16
   3.1. Introduction ................................................................................................................ 16
   3.2. Views .......................................................................................................................... 16
   3.3. Foreign Keys ................................................................................................................ 16
   3.4. Transactions ............................................................................................................... 17
   3.5. Window Functions ..................................................................................................... 19
   3.6. Inheritance ............................................................................................................... 22
   3.7. Conclusion .................................................................................................................. 23
Chapter 1. Getting Started

1.1. Installation

Before you can use PostgreSQL you need to install it, of course. It is possible that PostgreSQL is already installed at your site, either because it was included in your operating system distribution or because the system administrator already installed it. If that is the case, you should obtain information from the operating system documentation or your system administrator about how to access PostgreSQL.

If you are not sure whether PostgreSQL is already available or whether you can use it for your experimentation then you can install it yourself. Doing so is not hard and it can be a good exercise. PostgreSQL can be installed by any unprivileged user; no superuser (root) access is required.

If you are installing PostgreSQL yourself, then refer to Chapter 16 for instructions on installation, and return to this guide when the installation is complete. Be sure to follow closely the section about setting up the appropriate environment variables.

If your site administrator has not set things up in the default way, you might have some more work to do. For example, if the database server machine is a remote machine, you will need to set the \texttt{PGHOST} environment variable to the name of the database server machine. The environment variable \texttt{PGPORT} might also have to be set. The bottom line is this: if you try to start an application program and it complains that it cannot connect to the database, you should consult your site administrator or, if that is you, the documentation to make sure that your environment is properly set up. If you did not understand the preceding paragraph then read the next section.

1.2. Architectural Fundamentals

Before we proceed, you should understand the basic PostgreSQL system architecture. Understanding how the parts of PostgreSQL interact will make this chapter somewhat clearer.

In database jargon, PostgreSQL uses a client/server model. A PostgreSQL session consists of the following cooperating processes (programs):

- A server process, which manages the database files, accepts connections to the database from client applications, and performs database actions on behalf of the clients. The database server program is called \texttt{postgres}.

- The user's client (frontend) application that wants to perform database operations. Client applications can be very diverse in nature: a client could be a text-oriented tool, a graphical application, a web server that accesses the database to display web pages, or a specialized database maintenance tool. Some client applications are supplied with the PostgreSQL distribution; most are developed by users.

As is typical of client/server applications, the client and the server can be on different hosts. In that case they communicate over a TCP/IP network connection. You should keep this in mind, because the files that can be accessed on a client machine might not be accessible (or might only be accessible using a different file name) on the database server machine.

The PostgreSQL server can handle multiple concurrent connections from clients. To achieve this it starts ("forks") a new process for each connection. From that point on, the client and the new server process communicate without intervention by the original \texttt{postgres} process. Thus, the master server process is always running, waiting for client connections, whereas client and associated server processes come and go. (All of this is of course invisible to the user. We only mention it here for completeness.)

1.3. Creating a Database
The first test to see whether you can access the database server is to try to create a database. A running PostgreSQL server can manage many databases. Typically, a separate database is used for each project or for each user.

Possibly, your site administrator has already created a database for your use. In that case you can omit this step and skip ahead to the next section.

To create a new database, in this example named mydb, you use the following command:

\[$ \text{createdb mydb} \$

If this produces no response then this step was successful and you can skip over the remainder of this section.

If you see a message similar to:

\[ \text{createdb: command not found} \]

then PostgreSQL was not installed properly. Either it was not installed at all or your shell's search path was not set to include it. Try calling the command with an absolute path instead:

\[$ \text{/usr/local/pgsql/bin/createdb mydb} \$

The path at your site might be different. Contact your site administrator or check the installation instructions to correct the situation.

Another response could be this:

\[ \text{createdb: could not connect to database postgres: could not connect to server: No such file or directory} \]
\[ \text{Is the server running locally and accepting connections on Unix domain socket "/tmp/.s.PGSQL.5432"?} \]

This means that the server was not started, or it was not started where createdb expected it. Again, check the installation instructions or consult the administrator.

Another response could be this:

\[ \text{createdb: could not connect to database postgres: FATAL: role "joe" does not exist} \]

where your own login name is mentioned. This will happen if the administrator has not created a PostgreSQL user account for you. (PostgreSQL user accounts are distinct from operating system user accounts.) If you are the administrator, see Chapter 21 for help creating accounts. You will need to become the operating system user under which PostgreSQL was installed (usually postgres) to create the first user account. It could also be that you were assigned a PostgreSQL user name that is different from your operating system user name; in that case you need to use the \(-U\) switch or set the \text{PGUSER} environment variable to specify your PostgreSQL user name.

If you have a user account but it does not have the privileges required to create a database, you will see the following:

\[ \text{createdb: database creation failed: ERROR: permission denied to create database} \]
Not every user has authorization to create new databases. If PostgreSQL refuses to create databases for you then the site administrator needs to grant you permission to create databases. Consult your site administrator if this occurs. If you installed PostgreSQL yourself then you should log in for the purposes of this tutorial under the user account that you started the server as. ¹

You can also create databases with other names. PostgreSQL allows you to create any number of databases at a given site. Database names must have an alphabetic first character and are limited to 63 bytes in length. A convenient choice is to create a database with the same name as your current user name. Many tools assume that database name as the default, so it can save you some typing. To create that database, simply type:

```
$ createdb
```

If you do not want to use your database anymore you can remove it. For example, if you are the owner (creator) of the database `mydb`, you can destroy it using the following command:

```
$ dropdb mydb
```

(For this command, the database name does not default to the user account name. You always need to specify it.) This action physically removes all files associated with the database and cannot be undone, so this should only be done with a great deal of forethought.

More about `createdb` and `dropdb` can be found in `createdb` and `dropdb` respectively.

### 1.4. Accessing a Database

Once you have created a database, you can access it by:

- Running the PostgreSQL interactive terminal program, called `psql`, which allows you to interactively enter, edit, and execute SQL commands.
- Using an existing graphical frontend tool like pgAdmin or an office suite with ODBC or JDBC support to create and manipulate a database. These possibilities are not covered in this tutorial.
- Writing a custom application, using one of the several available language bindings. These possibilities are discussed further in Part IV.

You probably want to start up `psql` to try the examples in this tutorial. It can be activated for the `mydb` database by typing the command:

```
$ psql mydb
```

If you do not supply the database name then it will default to your user account name. You already discovered this scheme in the previous section using `createdb`.

In `psql`, you will be greeted with the following message:

```
psql (11.1)
Type "help" for help.

mydb=>
```

The last line could also be:

---

¹ As an explanation for why this works: PostgreSQL user names are separate from operating system user accounts. When you connect to a database, you can choose what PostgreSQL user name to connect as; if you don’t, it will default to the same name as your current operating system account. As it happens, there will always be a PostgreSQL user account that has the same name as the operating system user that started the server, and it also happens that that user always has permission to create databases. Instead of logging in as that user you can also specify the `-U` option everywhere to select a PostgreSQL user name to connect as.
That would mean you are a database superuser, which is most likely the case if you installed the PostgreSQL instance yourself. Being a superuser means that you are not subject to access controls. For the purposes of this tutorial that is not important.

If you encounter problems starting psql then go back to the previous section. The diagnostics of createdb and psql are similar, and if the former worked the latter should work as well.

The last line printed out by psql is the prompt, and it indicates that psql is listening to you and that you can type SQL queries into a work space maintained by psql. Try out these commands:

```
mydb=> SELECT version();
version
------------------------------------------------------------------------------------------
 PostgreSQL 11.1 on x86_64-pc-linux-gnu, compiled by gcc (Debian 4.9.2-10) 4.9.2, 64-bit
(1 row)

mydb=> SELECT current_date;
date
------------
2016-01-07
(1 row)

mydb=> SELECT 2 + 2;
?column?
----------
4
(1 row)
```

The psql program has a number of internal commands that are not SQL commands. They begin with the backslash character, "\". For example, you can get help on the syntax of various PostgreSQL SQL commands by typing:

```
mydb=> \h
```

To get out of psql, type:

```
mydb=> \q
```

and psql will quit and return you to your command shell. (For more internal commands, type \? at the psql prompt.) The full capabilities of psql are documented in psql. In this tutorial we will not use these features explicitly, but you can use them yourself when it is helpful.
Chapter 2. The SQL Language

2.1. Introduction

This chapter provides an overview of how to use SQL to perform simple operations. This tutorial is only intended to give you an introduction and is in no way a complete tutorial on SQL. Numerous books have been written on SQL, including [mel93] and [date97]. You should be aware that some PostgreSQL language features are extensions to the standard.

In the examples that follow, we assume that you have created a database named mydb, as described in the previous chapter, and have been able to start psql.

Examples in this manual can also be found in the PostgreSQL source distribution in the directory src/tutorial/. (Binary distributions of PostgreSQL might not compile these files.) To use those files, first change to that directory and run make:

$$
cd ....~/src/tutorial
make
$$

This creates the scripts and compiles the C files containing user-defined functions and types. Then, to start the tutorial, do the following:

$$
cd ....~/tutorial
psql -s mydb
... 
mydb=> \i basics.sql
$$

The \i command reads in commands from the specified file. psql's -s option puts you in single step mode which pauses before sending each statement to the server. The commands used in this section are in the file basics.sql.

2.2. Concepts

PostgreSQL is a relational database management system (RDBMS). That means it is a system for managing data stored in relations. Relation is essentially a mathematical term for table. The notion of storing data in tables is so commonplace today that it might seem inherently obvious, but there are a number of other ways of organizing databases. Files and directories on Unix-like operating systems form an example of a hierarchical database. A more modern development is the object-oriented database.

Each table is a named collection of rows. Each row of a given table has the same set of named columns, and each column is of a specific data type. Whereas columns have a fixed order in each row, it is important to remember that SQL does not guarantee the order of the rows within the table in any way (although they can be explicitly sorted for display).

Tables are grouped into databases, and a collection of databases managed by a single PostgreSQL server instance constitutes a database cluster.

2.3. Creating a New Table

You can create a new table by specifying the table name, along with all column names and their types:
The SQL Language

CREATE TABLE weather (  
city varchar(80),  
temp_lo int, -- low temperature  
temp_hi int, -- high temperature  
prcp real, -- precipitation  
date date
);

You can enter this into psql with the line breaks. psql will recognize that the command is not terminated until the semicolon.

White space (i.e., spaces, tabs, and newlines) can be used freely in SQL commands. That means you can type the command aligned differently than above, or even all on one line. Two dashes ("--") introduce comments. Whatever follows them is ignored up to the end of the line. SQL is case insensitive about key words and identifiers, except when identifiers are double-quoted to preserve the case (not done above).

varchar(80) specifies a data type that can store arbitrary character strings up to 80 characters in length. int is the normal integer type. real is a type for storing single precision floating-point numbers. date should be self-explanatory. (Yes, the column of type date is also named date. This might be convenient or confusing — you choose.)

PostgreSQL supports the standard SQL types int, smallint, real, double precision, char(N), varchar(N), date, time, timestamp, and interval, as well as other types of general utility and a rich set of geometric types. PostgreSQL can be customized with an arbitrary number of user-defined data types. Consequently, type names are not key words in the syntax, except where required to support special cases in the SQL standard.

The second example will store cities and their associated geographical location:

CREATE TABLE cities (  
name varchar(80),  
location point
);

The point type is an example of a PostgreSQL-specific data type.

Finally, it should be mentioned that if you don’t need a table any longer or want to recreate it differently you can remove it using the following command:

DROP TABLE tablename;

2.4. Populating a Table With Rows

The INSERT statement is used to populate a table with rows:

INSERT INTO weather VALUES ('San Francisco', 46, 50, 0.25, '1994-11-27');

Note that all data types use rather obvious input formats. Constants that are not simple numeric values usually must be surrounded by single quotes (""), as in the example. The date type is actually quite flexible in what it accepts, but for this tutorial we will stick to the unambiguous format shown here.

The point type requires a coordinate pair as input, as shown here:
The SQL Language

```
INSERT INTO cities VALUES ('San Francisco', '(-194.0, 53.0)');
```

The syntax used so far requires you to remember the order of the columns. An alternative syntax allows you to list the columns explicitly:

```
INSERT INTO weather (city, temp_lo, temp_hi, prcp, date)
VALUES ('San Francisco', 43, 57, 0.0, '1994-11-29');
```

You can list the columns in a different order if you wish or even omit some columns, e.g., if the precipitation is unknown:

```
INSERT INTO weather (date, city, temp_hi, temp_lo)
VALUES ('1994-11-29', 'Hayward', 54, 37);
```

Many developers consider explicitly listing the columns better style than relying on the order implicitly.

Please enter all the commands shown above so you have some data to work with in the following sections.

You could also have used `COPY` to load large amounts of data from flat-text files. This is usually faster because the `COPY` command is optimized for this application while allowing less flexibility than `INSERT`. An example would be:

```
COPY weather FROM '/home/user/weather.txt';
```

where the file name for the source file must be available on the machine running the backend process, not the client, since the backend process reads the file directly. You can read more about the `COPY` command in `COPY`.

### 2.5. Querying a Table

To retrieve data from a table, the table is *queried*. An SQL `SELECT` statement is used to do this. The statement is divided into a select list (the part that lists the columns to be returned), a table list (the part that lists the tables from which to retrieve the data), and an optional qualification (the part that specifies any restrictions). For example, to retrieve all the rows of table `weather`, type:

```
SELECT * FROM weather;
```

Here `*` is a shorthand for "all columns". So the same result would be had with:

```
SELECT city, temp_lo, temp_hi, prcp, date FROM weather;
```

The output should be:

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>43</td>
<td>57</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>Hayward</td>
<td>37</td>
<td>54</td>
<td></td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>

(3 rows)

---

1 While `SELECT *` is useful for off-the-cuff queries, it is widely considered bad style in production code, since adding a column to the table would change the results.
You can write expressions, not just simple column references, in the select list. For example, you can do:

```
SELECT city, (temp_hi+temp_lo)/2 AS temp_avg, date FROM weather;
```

This should give:

<table>
<thead>
<tr>
<th>city</th>
<th>temp_avg</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>48</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>50</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>Hayward</td>
<td>45</td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>

(3 rows)

Notice how the AS clause is used to relabel the output column. (The AS clause is optional.)

A query can be “qualified” by adding a WHERE clause that specifies which rows are wanted. The WHERE clause contains a Boolean (truth value) expression, and only rows for which the Boolean expression is true are returned. The usual Boolean operators (AND, OR, and NOT) are allowed in the qualification. For example, the following retrieves the weather of San Francisco on rainy days:

```
SELECT * FROM weather
WHERE city = 'San Francisco' AND prcp > 0.0;
```

Result:

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
</tbody>
</table>

(1 row)

You can request that the results of a query be returned in sorted order:

```
SELECT * FROM weather
ORDER BY city;
```

In this example, the sort order isn't fully specified, and so you might get the San Francisco rows in either order. But you'd always get the results shown above if you do:

```
SELECT * FROM weather
ORDER BY city, temp_lo;
```

You can request that duplicate rows be removed from the result of a query:

```
SELECT DISTINCT city
FROM weather;
```


Here again, the result row ordering might vary. You can ensure consistent results by using `DISTINCT` and `ORDER BY` together:

```sql
SELECT DISTINCT city
FROM weather
ORDER BY city;
```

### 2.6. Joins Between Tables

Thus far, our queries have only accessed one table at a time. Queries can access multiple tables at once, or access the same table in such a way that multiple rows of the table are being processed at the same time. A query that accesses multiple rows of the same or different tables at one time is called a join query. As an example, say you wish to list all the weather records together with the location of the associated city. To do that, we need to compare the `city` column of each row of the `weather` table with the `name` column of all rows in the `cities` table, and select the pairs of rows where these values match.

This would be accomplished by the following query:

```sql
SELECT *
FROM weather, cities
WHERE city = name;
```

#### Note

This is only a conceptual model. The join is usually performed in a more efficient manner than actually comparing each possible pair of rows, but this is invisible to the user.

Observe two things about the result set:

- There is no result row for the city of Hayward. This is because there is no matching entry in the `cities` table for Hayward, so the join ignores the unmatched rows in the `weather` table. We will see shortly how this can be fixed.

---

2 In some database systems, including older versions of PostgreSQL, the implementation of `DISTINCT` automatically orders the rows and `ORDER BY` is unnecessary. But this is not required by the SQL standard, and current PostgreSQL does not guarantee that `DISTINCT` causes the rows to be ordered.
• There are two columns containing the city name. This is correct because the lists of columns from
the weather and cities tables are concatenated. In practice this is undesirable, though, so you
will probably want to list the output columns explicitly rather than using *:

```
SELECT city, temp_lo, temp_hi, prcp, date, location
FROM weather, cities
WHERE city = name;
```

**Exercise:** Attempt to determine the semantics of this query when the WHERE clause is omitted.

Since the columns all had different names, the parser automatically found which table they belong
to. If there were duplicate column names in the two tables you’d need to qualify the column names
to show which one you meant, as in:

```
SELECT weather.city, weather.temp_lo, weather.temp_hi,
    weather.prcp, weather.date, cities.location
FROM weather, cities
WHERE cities.name = weather.city;
```

It is widely considered good style to qualify all column names in a join query, so that the query won’t
fail if a duplicate column name is later added to one of the tables.

Join queries of the kind seen thus far can also be written in this alternative form:

```
SELECT *
    FROM weather INNER JOIN cities ON (weather.city = cities.name);
```

This syntax is not as commonly used as the one above, but we show it here to help you understand
the following topics.

Now we will figure out how we can get the Hayward records back in. What we want the query to do
is to scan the weather table and for each row to find the matching cities row(s). If no matching
row is found we want some “empty values” to be substituted for the cities table’s columns. This
kind of query is called an **outer join**. (The joins we have seen so far are inner joins.) The command
looks like this:

```
SELECT *
    FROM weather LEFT OUTER JOIN cities ON (weather.city = cities.name);
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
<th>name</th>
<th>location</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hayward</td>
<td>37</td>
<td>54</td>
<td></td>
<td>1994-11-29</td>
<td></td>
<td></td>
</tr>
<tr>
<td>San</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
<td>San Francisco</td>
<td></td>
</tr>
<tr>
<td>Francisco</td>
<td>(-194,53)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>San</td>
<td>43</td>
<td>57</td>
<td></td>
<td>1994-11-29</td>
<td>San Francisco</td>
<td></td>
</tr>
<tr>
<td>Francisco</td>
<td>(-194,53)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

(3 rows)

This query is called a **left outer join** because the table mentioned on the left of the join operator will
have each of its rows in the output at least once, whereas the table on the right will only have those
rows output that match some row of the left table. When outputting a left-table row for which there is
no right-table match, empty (null) values are substituted for the right-table columns.
**Exercise:** There are also right outer joins and full outer joins. Try to find out what those do.

We can also join a table against itself. This is called a *self join*. As an example, suppose we wish to find all the weather records that are in the temperature range of other weather records. So we need to compare the `temp_lo` and `temp_hi` columns of each weather row to the `temp_lo` and `temp_hi` columns of all other weather rows. We can do this with the following query:

```sql
SELECT W1.city, W1.temp_lo AS low, W1.temp_hi AS high,
       W2.city, W2.temp_lo AS low, W2.temp_hi AS high
FROM weather W1, weather W2
WHERE W1.temp_lo < W2.temp_lo
  AND W1.temp_hi > W2.temp_hi;
```

<table>
<thead>
<tr>
<th>city</th>
<th>low</th>
<th>high</th>
<th>city</th>
<th>low</th>
<th>high</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>43</td>
<td>57</td>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
</tr>
<tr>
<td>Hayward</td>
<td>37</td>
<td>54</td>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
</tr>
</tbody>
</table>

(2 rows)

Here we have relabeled the weather table as `W1` and `W2` to be able to distinguish the left and right side of the join. You can also use these kinds of aliases in other queries to save some typing, e.g.:

```sql
SELECT *
FROM weather w, cities c
WHERE w.city = c.name;
```

You will encounter this style of abbreviating quite frequently.

### 2.7. Aggregate Functions

Like most other relational database products, PostgreSQL supports *aggregate functions*. An aggregate function computes a single result from multiple input rows. For example, there are aggregates to compute the `count`, `sum`, `avg` (average), `max` (maximum) and `min` (minimum) over a set of rows.

As an example, we can find the highest low-temperature reading anywhere with:

```sql
SELECT max(temp_lo) FROM weather;
```

```
max
-----
46
(1 row)
```

If we wanted to know what city (or cities) that reading occurred in, we might try:

```sql
SELECT city FROM weather WHERE temp_lo = max(temp_lo);
```

but this will not work since the aggregate `max` cannot be used in the `WHERE` clause. (This restriction exists because the `WHERE` clause determines which rows will be included in the aggregate calculation; so obviously it has to be evaluated before aggregate functions are computed.) However, as is often the case the query can be restated to accomplish the desired result, here by using a *subquery*:

```sql
SELECT city FROM weather
```
WHERE temp_lo = (SELECT max(temp_lo) FROM weather);

city
-------------
San Francisco
(1 row)

This is OK because the subquery is an independent computation that computes its own aggregate separately from what is happening in the outer query.

Aggregates are also very useful in combination with GROUP BY clauses. For example, we can get the maximum low temperature observed in each city with:

```
SELECT city, max(temp_lo)
FROM weather
GROUP BY city;
```

city      | max
-------------+-----
Hayward       |  37
San Francisco |  46
(2 rows)

which gives us one output row per city. Each aggregate result is computed over the table rows matching that city. We can filter these grouped rows using HAVING:

```
SELECT city, max(temp_lo)
FROM weather
GROUP BY city
HAVING max(temp_lo) < 40;
```

city   | max
---------+-----
Hayward |  37
(1 row)

which gives us the same results for only the cities that have all temp_lo values below 40. Finally, if we only care about cities whose names begin with “S”, we might do:

```
SELECT city, max(temp_lo)
FROM weather
WHERE city LIKE 'S%'            -- 1
GROUP BY city
HAVING max(temp_lo) < 40;
```

1 The LIKE operator does pattern matching and is explained in Section 9.7.

It is important to understand the interaction between aggregates and SQL’s WHERE and HAVING clauses. The fundamental difference between WHERE and HAVING is this: WHERE selects input rows before groups and aggregates are computed (thus, it controls which rows go into the aggregate computation), whereas HAVING selects group rows after groups and aggregates are computed. Thus, the WHERE clause must not contain aggregate functions; it makes no sense to try to use an aggregate to determine which rows will be inputs to the aggregates. On the other hand, the HAVING clause always contains aggregate functions. (Strictly speaking, you are allowed to write a HAVING clause that doesn’t use
aggregates, but it’s seldom useful. The same condition could be used more efficiently at the WHERE stage.)

In the previous example, we can apply the city name restriction in WHERE, since it needs no aggregate. This is more efficient than adding the restriction to HAVING, because we avoid doing the grouping and aggregate calculations for all rows that fail the WHERE check.

2.8. Updates

You can update existing rows using the UPDATE command. Suppose you discover the temperature readings are all off by 2 degrees after November 28. You can correct the data as follows:

```
UPDATE weather
SET temp_hi = temp_hi - 2, temp_lo = temp_lo - 2
WHERE date > '1994-11-28';
```

Look at the new state of the data:

```
SELECT * FROM weather;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>41</td>
<td>55</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>Hayward</td>
<td>35</td>
<td>52</td>
<td></td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>

(3 rows)

2.9. Deletions

Rows can be removed from a table using the DELETE command. Suppose you are no longer interested in the weather of Hayward. Then you can do the following to delete those rows from the table:

```
DELETE FROM weather WHERE city = 'Hayward';
```

All weather records belonging to Hayward are removed.

```
SELECT * FROM weather;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>41</td>
<td>55</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>

(2 rows)

One should be wary of statements of the form

```
DELETE FROM tablename;
```

Without a qualification, DELETE will remove all rows from the given table, leaving it empty. The system will not request confirmation before doing this!
Chapter 3. Advanced Features

3.1. Introduction

In the previous chapter we have covered the basics of using SQL to store and access your data in PostgreSQL. We will now discuss some more advanced features of SQL that simplify management and prevent loss or corruption of your data. Finally, we will look at some PostgreSQL extensions.

This chapter will on occasion refer to examples found in Chapter 2 to change or improve them, so it will be useful to have read that chapter. Some examples from this chapter can also be found in advanced.sql in the tutorial directory. This file also contains some sample data to load, which is not repeated here. (Refer to Section 2.1 for how to use the file.)

3.2. Views

Refer back to the queries in Section 2.6. Suppose the combined listing of weather records and city location is of particular interest to your application, but you do not want to type the query each time you need it. You can create a view over the query, which gives a name to the query that you can refer to like an ordinary table:

```
CREATE VIEW myview AS
    SELECT city, temp_lo, temp_hi, prcp, date, location
    FROM weather, cities
    WHERE city = name;

SELECT * FROM myview;
```

Making liberal use of views is a key aspect of good SQL database design. Views allow you to encapsulate the details of the structure of your tables, which might change as your application evolves, behind consistent interfaces.

Views can be used in almost any place a real table can be used. Building views upon other views is not uncommon.

3.3. Foreign Keys

Recall the weather and cities tables from Chapter 2. Consider the following problem: You want to make sure that no one can insert rows in the weather table that do not have a matching entry in the cities table. This is called maintaining the referential integrity of your data. In simplistic database systems this would be implemented (if at all) by first looking at the cities table to check if a matching record exists, and then inserting or rejecting the new weather records. This approach has a number of problems and is very inconvenient, so PostgreSQL can do this for you.

The new declaration of the tables would look like this:

```
CREATE TABLE cities (    
    city varchar(80) primary key,    
    location point    
));

CREATE TABLE weather (    
    city varchar(80) references cities(city),    
    temp_lo int,
```
Now try inserting an invalid record:

```
INSERT INTO weather VALUES ('Berkeley', 45, 53, 0.0, '1994-11-28');
```

ERROR:  insert or update on table "weather" violates foreign key constraint "weather_city_fkey"
DETAIL:  Key (city)=(Berkeley) is not present in table "cities".

The behavior of foreign keys can be finely tuned to your application. We will not go beyond this simple example in this tutorial, but just refer you to Chapter 5 for more information. Making correct use of foreign keys will definitely improve the quality of your database applications, so you are strongly encouraged to learn about them.

### 3.4. Transactions

*Transactions* are a fundamental concept of all database systems. The essential point of a transaction is that it bundles multiple steps into a single, all-or-nothing operation. The intermediate states between the steps are not visible to other concurrent transactions, and if some failure occurs that prevents the transaction from completing, then none of the steps affect the database at all.

For example, consider a bank database that contains balances for various customer accounts, as well as total deposit balances for branches. Suppose that we want to record a payment of $100.00 from Alice's account to Bob's account. Simplifying outrageously, the SQL commands for this might look like:

```
UPDATE accounts SET balance = balance - 100.00
    WHERE name = 'Alice';
UPDATE branches SET balance = balance - 100.00
    WHERE name = (SELECT branch_name FROM accounts WHERE name = 'Alice');
UPDATE accounts SET balance = balance + 100.00
    WHERE name = 'Bob';
UPDATE branches SET balance = balance + 100.00
    WHERE name = (SELECT branch_name FROM accounts WHERE name = 'Bob');
```

The details of these commands are not important here; the important point is that there are several separate updates involved to accomplish this rather simple operation. Our bank's officers will want to be assured that either all these updates happen, or none of them happen. It would certainly not do for a system failure to result in Bob receiving $100.00 that was not debited from Alice. Nor would Alice long remain a happy customer if she was debited without Bob being credited. We need a guarantee that if something goes wrong partway through the operation, none of the steps executed so far will take effect. Grouping the updates into a *transaction* gives us this guarantee. A transaction is said to be *atomic*: from the point of view of other transactions, it either happens completely or not at all.

We also want a guarantee that once a transaction is completed and acknowledged by the database system, it has indeed been permanently recorded and won't be lost even if a crash ensues shortly thereafter. For example, if we are recording a cash withdrawal by Bob, we do not want any chance that the debit to his account will disappear in a crash just after he walks out the bank door. A transactional database guarantees that all the updates made by a transaction are logged in permanent storage (i.e., on disk) before the transaction is reported complete.
Another important property of transactional databases is closely related to the notion of atomic updates: when multiple transactions are running concurrently, each one should not be able to see the incomplete changes made by others. For example, if one transaction is busy totalling all the branch balances, it would not do for it to include the debit from Alice's branch but not the credit to Bob's branch, nor vice versa. So transactions must be all-or-nothing not only in terms of their permanent effect on the database, but also in terms of their visibility as they happen. The updates made so far by an open transaction are invisible to other transactions until the transaction completes, whereupon all the updates become visible simultaneously.

In PostgreSQL, a transaction is set up by surrounding the SQL commands of the transaction with BEGIN and COMMIT commands. So our banking transaction would actually look like:

```sql
BEGIN;
UPDATE accounts SET balance = balance - 100.00
  WHERE name = 'Alice';
-- etc etc
COMMIT;
```

If, partway through the transaction, we decide we do not want to commit (perhaps we just noticed that Alice's balance went negative), we can issue the command ROLLBACK instead of COMMIT, and all our updates so far will be canceled.

PostgreSQL actually treats every SQL statement as being executed within a transaction. If you do not issue a BEGIN command, then each individual statement has an implicit BEGIN and (if successful) COMMIT wrapped around it. A group of statements surrounded by BEGIN and COMMIT is sometimes called a **transaction block**.

---

**Note**

Some client libraries issue BEGIN and COMMIT commands automatically, so that you might get the effect of transaction blocks without asking. Check the documentation for the interface you are using.

---

It's possible to control the statements in a transaction in a more granular fashion through the use of savepoints. Savepoints allow you to selectively discard parts of the transaction, while committing the rest. After defining a savepoint with SAVEPOINT, you can if needed roll back to the savepoint with ROLLBACK TO. All the transaction's database changes between defining the savepoint and rolling back to it are discarded, but changes earlier than the savepoint are kept.

After rolling back to a savepoint, it continues to be defined, so you can roll back to it several times. Conversely, if you are sure you won't need to roll back to a particular savepoint again, it can be released, so the system can free some resources. Keep in mind that either releasing or rolling back to a savepoint will automatically release all savepoints that were defined after it.

All this is happening within the transaction block, so none of it is visible to other database sessions. When and if you commit the transaction block, the committed actions become visible as a unit to other sessions, while the rolled-back actions never become visible at all.

Remembering the bank database, suppose we debit $100.00 from Alice's account, and credit Bob's account, only to find later that we should have credited Wally's account. We could do it using savepoints like this:

```sql
BEGIN;
UPDATE accounts SET balance = balance - 100.00
  WHERE name = 'Alice';
SAVEPOINT my_savepoint;
```
UPDATE accounts SET balance = balance + 100.00
  WHERE name = 'Bob';
-- oops ... forget that and use Wally's account
ROLLBACK TO my_savepoint;
UPDATE accounts SET balance = balance + 100.00
  WHERE name = 'Wally';
COMMIT;

This example is, of course, oversimplified, but there's a lot of control possible in a transaction block through the use of savepoints. Moreover, ROLLBACK TO is the only way to regain control of a transaction block that was put in aborted state by the system due to an error, short of rolling it back completely and starting again.

### 3.5. Window Functions

A **window function** performs a calculation across a set of table rows that are somehow related to the current row. This is comparable to the type of calculation that can be done with an aggregate function. However, window functions do not cause rows to become grouped into a single output row like non-window aggregate calls would. Instead, the rows retain their separate identities. Behind the scenes, the window function is able to access more than just the current row of the query result.

Here is an example that shows how to compare each employee's salary with the average salary in his or her department:

```
SELECT depname, empno, salary, avg(salary) OVER (PARTITION BY depname) FROM empsalary;
```

<table>
<thead>
<tr>
<th>depname</th>
<th>empno</th>
<th>salary</th>
<th>avg</th>
</tr>
</thead>
<tbody>
<tr>
<td>develop</td>
<td>11</td>
<td>5200</td>
<td>5020.000000000000000000</td>
</tr>
<tr>
<td>develop</td>
<td>7</td>
<td>4200</td>
<td>5020.000000000000000000</td>
</tr>
<tr>
<td>develop</td>
<td>9</td>
<td>4500</td>
<td>5020.000000000000000000</td>
</tr>
<tr>
<td>develop</td>
<td>8</td>
<td>6000</td>
<td>5020.000000000000000000</td>
</tr>
<tr>
<td>develop</td>
<td>10</td>
<td>5200</td>
<td>5020.000000000000000000</td>
</tr>
<tr>
<td>personnel</td>
<td>5</td>
<td>3500</td>
<td>3700.000000000000000000</td>
</tr>
<tr>
<td>personnel</td>
<td>2</td>
<td>3900</td>
<td>3700.000000000000000000</td>
</tr>
<tr>
<td>sales</td>
<td>3</td>
<td>4800</td>
<td>4866.6666666666666667</td>
</tr>
<tr>
<td>sales</td>
<td>1</td>
<td>5000</td>
<td>4866.6666666666666667</td>
</tr>
<tr>
<td>sales</td>
<td>4</td>
<td>4800</td>
<td>4866.6666666666666667</td>
</tr>
</tbody>
</table>

(10 rows)

The first three output columns come directly from the table `empsalary`, and there is one output row for each row in the table. The fourth column represents an average taken across all the table rows that have the same `depname` value as the current row. (This actually is the same function as the non-window `avg` aggregate, but the `OVER` clause causes it to be treated as a window function and computed across the window frame.)

A window function call always contains an `OVER` clause directly following the window function's name and argument(s). This is what syntactically distinguishes it from a normal function or non-window aggregate. The `OVER` clause determines exactly how the rows of the query are split up for processing by the window function. The `PARTITION BY` clause within `OVER` divides the rows into groups, or partitions, that share the same values of the `PARTITION BY` expression(s). For each row, the window function is computed across the rows that fall into the same partition as the current row.

You can also control the order in which rows are processed by window functions using `ORDER BY` within `OVER`. (The window `ORDER BY` does not even have to match the order in which the rows are output.) Here is an example:
SELECT depname, empno, salary,
    rank() OVER (PARTITION BY depname ORDER BY salary DESC)
FROM empsalary;

<table>
<thead>
<tr>
<th>depname</th>
<th>empno</th>
<th>salary</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>develop</td>
<td>8</td>
<td>6000</td>
<td>1</td>
</tr>
<tr>
<td>develop</td>
<td>10</td>
<td>5200</td>
<td>2</td>
</tr>
<tr>
<td>develop</td>
<td>11</td>
<td>5200</td>
<td>2</td>
</tr>
<tr>
<td>develop</td>
<td>9</td>
<td>4500</td>
<td>4</td>
</tr>
<tr>
<td>develop</td>
<td>7</td>
<td>4200</td>
<td>5</td>
</tr>
<tr>
<td>personnel</td>
<td>2</td>
<td>3900</td>
<td>1</td>
</tr>
<tr>
<td>personnel</td>
<td>5</td>
<td>3500</td>
<td>2</td>
</tr>
<tr>
<td>sales</td>
<td>1</td>
<td>5000</td>
<td>1</td>
</tr>
<tr>
<td>sales</td>
<td>4</td>
<td>4800</td>
<td>2</td>
</tr>
<tr>
<td>sales</td>
<td>3</td>
<td>4800</td>
<td>2</td>
</tr>
</tbody>
</table>

(10 rows)

As shown here, the `rank` function produces a numerical rank for each distinct `ORDER BY` value in the current row's partition, using the order defined by the `ORDER BY` clause. `rank` needs no explicit parameter, because its behavior is entirely determined by the `OVER` clause.

The rows considered by a window function are those of the “virtual table” produced by the query’s `FROM` clause as filtered by its `WHERE`, `GROUP BY`, and `HAVING` clauses if any. For example, a row removed because it does not meet the `WHERE` condition is not seen by any window function. A query can contain multiple window functions that slice up the data in different ways using different `OVER` clauses, but they all act on the same collection of rows defined by this virtual table.

We already saw that `ORDER BY` can be omitted if the ordering of rows is not important. It is also possible to omit `PARTITION BY`, in which case there is a single partition containing all rows.

There is another important concept associated with window functions: for each row, there is a set of rows within its partition called its `window frame`. Some window functions act only on the rows of the window frame, rather than of the whole partition. By default, if `ORDER BY` is supplied then the frame consists of all rows from the start of the partition up through the current row, plus any following rows that are equal to the current row according to the `ORDER BY` clause. When `ORDER BY` is omitted the default frame consists of all rows in the partition. ¹ Here is an example using `sum`:

```
SELECT salary, sum(salary) OVER () FROM empsalary;
```

<table>
<thead>
<tr>
<th>salary</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>5200</td>
<td>47110</td>
</tr>
<tr>
<td>5000</td>
<td>47110</td>
</tr>
<tr>
<td>3500</td>
<td>47110</td>
</tr>
<tr>
<td>4800</td>
<td>47110</td>
</tr>
<tr>
<td>3900</td>
<td>47110</td>
</tr>
<tr>
<td>4200</td>
<td>47110</td>
</tr>
<tr>
<td>4500</td>
<td>47110</td>
</tr>
<tr>
<td>4800</td>
<td>47110</td>
</tr>
<tr>
<td>6000</td>
<td>47110</td>
</tr>
<tr>
<td>5200</td>
<td>47110</td>
</tr>
</tbody>
</table>

(10 rows)

¹ There are options to define the window frame in other ways, but this tutorial does not cover them. See Section 4.2.8 for details.
Above, since there is no ORDER BY in the OVER clause, the window frame is the same as the partition, which for lack of PARTITION BY is the whole table; in other words each sum is taken over the whole table and so we get the same result for each output row. But if we add an ORDER BY clause, we get very different results:

```sql
SELECT salary, sum(salary) OVER (ORDER BY salary) FROM empsalary;
```

<table>
<thead>
<tr>
<th>salary</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>3500</td>
<td>3500</td>
</tr>
<tr>
<td>3900</td>
<td>7400</td>
</tr>
<tr>
<td>4200</td>
<td>11600</td>
</tr>
<tr>
<td>4500</td>
<td>16100</td>
</tr>
<tr>
<td>4800</td>
<td>25700</td>
</tr>
<tr>
<td>4800</td>
<td>25700</td>
</tr>
<tr>
<td>5000</td>
<td>30700</td>
</tr>
<tr>
<td>5200</td>
<td>41100</td>
</tr>
<tr>
<td>5200</td>
<td>41100</td>
</tr>
<tr>
<td>6000</td>
<td>47100</td>
</tr>
</tbody>
</table>

(10 rows)

Here the sum is taken from the first (lowest) salary up through the current one, including any duplicates of the current one (notice the results for the duplicated salaries).

Window functions are permitted only in the SELECT list and the ORDER BY clause of the query. They are forbidden elsewhere, such as in GROUP BY, HAVING and WHERE clauses. This is because they logically execute after the processing of those clauses. Also, window functions execute after non-window aggregate functions. This means it is valid to include an aggregate function call in the arguments of a window function, but not vice versa.

If there is a need to filter or group rows after the window calculations are performed, you can use a sub-select. For example:

```sql
SELECT depname, empno, salary, enroll_date
FROM
  (SELECT depname, empno, salary, enroll_date,
   rank() OVER (PARTITION BY depname ORDER BY salary DESC, empno) AS pos
   FROM empsalary
  ) AS ss
WHERE pos < 3;
```

The above query only shows the rows from the inner query having rank less than 3.

When a query involves multiple window functions, it is possible to write out each one with a separate OVER clause, but this is duplicative and error-prone if the same windowing behavior is wanted for several functions. Instead, each windowing behavior can be named in a WINDOW clause and then referenced in OVER. For example:

```sql
SELECT sum(salary) OVER w, avg(salary) OVER w
FROM empsalary
WINDOW w AS (PARTITION BY depname ORDER BY salary DESC);
```

More details about window functions can be found in Section 4.2.8, Section 9.21, Section 7.2.5, and the SELECT reference page.
3.6. Inheritance

Inheritance is a concept from object-oriented databases. It opens up interesting new possibilities of database design.

Let's create two tables: A table cities and a table capitals. Naturally, capitals are also cities, so you want some way to show the capitals implicitly when you list all cities. If you're really clever you might invent some scheme like this:

```sql
CREATE TABLE capitals (  
    name       text,  
    population real,  
    altitude   int,   -- (in ft)  
    state      char(2) 
);
CREATE TABLE non_capitals (  
    name       text,  
    population real,  
    altitude   int     -- (in ft)  
);
CREATE VIEW cities AS  
    SELECT name, population, altitude FROM capitals  
    UNION  
    SELECT name, population, altitude FROM non_capitals;
```

This works OK as far as querying goes, but it gets ugly when you need to update several rows, for one thing.

A better solution is this:

```sql
CREATE TABLE cities (  
    name       text,  
    population real,  
    altitude   int     -- (in ft)  
);
CREATE TABLE capitals (  
    state      char(2)  
) INHERITS (cities);
```

In this case, a row of capitals inherits all columns (name, population, and altitude) from its parent, cities. The type of the column name is text, a native PostgreSQL type for variable length character strings. State capitals have an extra column, state, that shows their state. In PostgreSQL, a table can inherit from zero or more other tables.

For example, the following query finds the names of all cities, including state capitals, that are located at an altitude over 500 feet:

```sql
SELECT name, altitude  
    FROM cities  
    WHERE altitude > 500;
```

which returns:
<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

(3 rows)

On the other hand, the following query finds all the cities that are not state capitals and are situated at an altitude over 500 feet:

```
SELECT name, altitude
FROM ONLY cities
WHERE altitude > 500;
```

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
</tbody>
</table>

(2 rows)

Here the `ONLY` before `cities` indicates that the query should be run over only the `cities` table, and not tables below `cities` in the inheritance hierarchy. Many of the commands that we have already discussed — `SELECT`, `UPDATE`, and `DELETE` — support this `ONLY` notation.

**Note**

Although inheritance is frequently useful, it has not been integrated with unique constraints or foreign keys, which limits its usefulness. See Section 5.9 for more detail.

### 3.7. Conclusion

PostgreSQL has many features not touched upon in this tutorial introduction, which has been oriented toward newer users of SQL. These features are discussed in more detail in the remainder of this book.

If you feel you need more introductory material, please visit the PostgreSQL web site\(^2\) for links to more resources.

---
\(^2\) [https://www.postgresql.org](https://www.postgresql.org)
Part II. The SQL Language

This part describes the use of the SQL language in PostgreSQL. We start with describing the general syntax of SQL, then explain how to create the structures to hold data, how to populate the database, and how to query it. The middle part lists the available data types and functions for use in SQL commands. The rest treats several aspects that are important for tuning a database for optimal performance.

The information in this part is arranged so that a novice user can follow it start to end to gain a full understanding of the topics without having to refer forward too many times. The chapters are intended to be self-contained, so that advanced users can read the chapters individually as they choose. The information in this part is presented in a narrative fashion in topical units. Readers looking for a complete description of a particular command should see Part VI.

Readers of this part should know how to connect to a PostgreSQL database and issue SQL commands. Readers that are unfamiliar with these issues are encouraged to read Part I first. SQL commands are typically entered using the PostgreSQL interactive terminal psql, but other programs that have similar functionality can be used as well.
# Table of Contents

4. SQL Syntax .................................................................................................................. 32
   4.1. Lexical Structure ..................................................................................................... 32
      4.1.1. Identifiers and Key Words .......................................................................... 32
      4.1.2. Constants ...................................................................................................... 34
      4.1.3. Operators ...................................................................................................... 38
      4.1.4. Special Characters ....................................................................................... 39
      4.1.5. Comments ..................................................................................................... 39
      4.1.6. Operator Precedence .................................................................................... 40
   4.2. Value Expressions .................................................................................................. 41
      4.2.1. Column References ....................................................................................... 42
      4.2.2. Positional Parameters ................................................................................... 42
      4.2.3. Subscripts ....................................................................................................... 42
      4.2.4. Field Selection ............................................................................................... 43
      4.2.5. Operator Invocations ..................................................................................... 43
      4.2.6. Function Calls ............................................................................................... 44
      4.2.7. Aggregate Expressions ................................................................................ 44
      4.2.8. Window Function Calls ............................................................................... 46
      4.2.9. Type Casts ..................................................................................................... 49
      4.2.10. Collation Expressions ................................................................................ 50
      4.2.11. Scalar Subqueries ....................................................................................... 51
      4.2.12. Array Constructors ................................................................................... 51
      4.2.13. Row Constructors ...................................................................................... 52
      4.2.14. Expression Evaluation Rules ...................................................................... 54
   4.3. Calling Functions ................................................................................................... 55
      4.3.1. Using Positional Notation ............................................................................ 56
      4.3.2. Using Named Notation ................................................................................ 56
      4.3.3. Using Mixed Notation ................................................................................ 57
5. Data Definition ............................................................................................................. 58
   5.1. Table Basics .......................................................................................................... 58
   5.2. Default Values ...................................................................................................... 59
   5.3. Constraints ........................................................................................................... 60
      5.3.1. Check Constraints ...................................................................................... 60
      5.3.2. Not-Null Constraints .................................................................................. 62
      5.3.3. Unique Constraints ..................................................................................... 63
      5.3.4. Primary Keys ............................................................................................... 63
      5.3.5. Foreign Keys ............................................................................................... 64
      5.3.6. Exclusion Constraints ................................................................................ 67
   5.4. System Columns .................................................................................................... 67
   5.5. Modifying Tables .................................................................................................. 68
      5.5.1. Adding a Column ......................................................................................... 69
      5.5.2. Removing a Column .................................................................................... 69
      5.5.3. Adding a Constraint .................................................................................... 69
      5.5.4. Removing a Constraint ............................................................................... 70
      5.5.5. Changing a Column's Default Value .......................................................... 70
      5.5.6. Changing a Column's Data Type ............................................................... 70
      5.5.7. Renaming a Column ................................................................................... 71
      5.5.8. Renaming a Table ....................................................................................... 71
   5.6. Privileges ................................................................................................................ 71
   5.7. Row Security Policies ............................................................................................ 72
   5.8. Schemas ................................................................................................................ 78
      5.8.1. Creating a Schema ...................................................................................... 79
      5.8.2. The Public Schema ...................................................................................... 79
      5.8.3. The Schema Search Path .......................................................................... 80
      5.8.4. Schemas and Privileges ............................................................................. 81
      5.8.5. The System Catalog Schema ..................................................................... 81
<table>
<thead>
<tr>
<th>Section Number</th>
<th>Section Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>8.13</td>
<td>XML Type</td>
<td>163</td>
</tr>
<tr>
<td>8.8</td>
<td>Geometric Types</td>
<td>154</td>
</tr>
<tr>
<td>8.8.1</td>
<td>Points</td>
<td>154</td>
</tr>
<tr>
<td>8.8.2</td>
<td>Lines</td>
<td>155</td>
</tr>
<tr>
<td>8.8.3</td>
<td>Line Segments</td>
<td>155</td>
</tr>
<tr>
<td>8.8.4</td>
<td>Boxes</td>
<td>155</td>
</tr>
<tr>
<td>8.8.5</td>
<td>Paths</td>
<td>156</td>
</tr>
<tr>
<td>8.8.6</td>
<td>Polygons</td>
<td>156</td>
</tr>
<tr>
<td>8.8.7</td>
<td>Circles</td>
<td>156</td>
</tr>
<tr>
<td>8.9</td>
<td>Network Address Types</td>
<td>156</td>
</tr>
<tr>
<td>8.9.1</td>
<td>inet</td>
<td>157</td>
</tr>
<tr>
<td>8.9.2</td>
<td>cidr</td>
<td>157</td>
</tr>
<tr>
<td>8.9.3</td>
<td>inet vs. cidr</td>
<td>158</td>
</tr>
<tr>
<td>8.9.4</td>
<td>macaddr</td>
<td>158</td>
</tr>
<tr>
<td>8.9.5</td>
<td>macaddr8</td>
<td>158</td>
</tr>
<tr>
<td>8.10</td>
<td>Bit String Types</td>
<td>159</td>
</tr>
<tr>
<td>8.11</td>
<td>Text Search Types</td>
<td>160</td>
</tr>
<tr>
<td>8.11.1</td>
<td>tsvector</td>
<td>160</td>
</tr>
<tr>
<td>8.11.2</td>
<td>tsquery</td>
<td>161</td>
</tr>
<tr>
<td>8.12</td>
<td>UUID Type</td>
<td>162</td>
</tr>
<tr>
<td>8.13.1</td>
<td>Creating XML Values</td>
<td>163</td>
</tr>
<tr>
<td>8.13.2</td>
<td>Encoding Handling</td>
<td>164</td>
</tr>
<tr>
<td>8.13.3</td>
<td>Accessing XML Values</td>
<td>165</td>
</tr>
<tr>
<td>8.14</td>
<td>JSON Types</td>
<td>165</td>
</tr>
<tr>
<td>8.14.1</td>
<td>JSON Input and Output Syntax</td>
<td>166</td>
</tr>
<tr>
<td>8.14.2</td>
<td>Designing JSON documents effectively</td>
<td>167</td>
</tr>
<tr>
<td>8.14.3</td>
<td>jsonb Containment and Existence</td>
<td>168</td>
</tr>
<tr>
<td>8.14.4</td>
<td>jsonb Indexing</td>
<td>170</td>
</tr>
<tr>
<td>8.14.5</td>
<td>Transforms</td>
<td>172</td>
</tr>
<tr>
<td>8.15</td>
<td>Arrays</td>
<td>172</td>
</tr>
<tr>
<td>8.15.1</td>
<td>Declaration of Array Types</td>
<td>172</td>
</tr>
<tr>
<td>8.15.2</td>
<td>Array Value Input</td>
<td>173</td>
</tr>
<tr>
<td>8.15.3</td>
<td>Accessing Arrays</td>
<td>174</td>
</tr>
<tr>
<td>8.15.4</td>
<td>Modifying Arrays</td>
<td>176</td>
</tr>
<tr>
<td>8.15.5</td>
<td>Searching in Arrays</td>
<td>179</td>
</tr>
<tr>
<td>8.15.6</td>
<td>Array Input and Output Syntax</td>
<td>180</td>
</tr>
<tr>
<td>8.16</td>
<td>Composite Types</td>
<td>181</td>
</tr>
<tr>
<td>8.16.1</td>
<td>Declaration of Composite Types</td>
<td>182</td>
</tr>
<tr>
<td>8.16.2</td>
<td>Constructing Composite Values</td>
<td>183</td>
</tr>
<tr>
<td>8.16.3</td>
<td>Accessing Composite Types</td>
<td>183</td>
</tr>
<tr>
<td>8.16.4</td>
<td>Modifying Composite Types</td>
<td>184</td>
</tr>
<tr>
<td>8.16.5</td>
<td>Using Composite Types in Queries</td>
<td>184</td>
</tr>
<tr>
<td>8.16.6</td>
<td>Composite Type Input and Output Syntax</td>
<td>187</td>
</tr>
<tr>
<td>8.17</td>
<td>Range Types</td>
<td>188</td>
</tr>
<tr>
<td>8.17.1</td>
<td>Built-in Range Types</td>
<td>188</td>
</tr>
<tr>
<td>8.17.2</td>
<td>Examples</td>
<td>188</td>
</tr>
<tr>
<td>8.17.3</td>
<td>Inclusive and Exclusive Bounds</td>
<td>189</td>
</tr>
<tr>
<td>8.17.4</td>
<td>Infinite (Unbounded) Ranges</td>
<td>189</td>
</tr>
<tr>
<td>8.17.5</td>
<td>Range Input/Output</td>
<td>189</td>
</tr>
<tr>
<td>8.17.6</td>
<td>Constructing Ranges</td>
<td>190</td>
</tr>
<tr>
<td>8.17.7</td>
<td>Discrete Range Types</td>
<td>191</td>
</tr>
<tr>
<td>8.17.8</td>
<td>Defining New Range Types</td>
<td>191</td>
</tr>
<tr>
<td>8.17.9</td>
<td>Indexing</td>
<td>192</td>
</tr>
<tr>
<td>8.17.10</td>
<td>Constraints on Ranges</td>
<td>192</td>
</tr>
<tr>
<td>8.18</td>
<td>Domain Types</td>
<td>193</td>
</tr>
<tr>
<td>8.19</td>
<td>Object Identifier Types</td>
<td>194</td>
</tr>
</tbody>
</table>
The SQL Language

8.20. pg_json Type ................................................................. 196
8.21. Pseudo-Types .............................................................. 196
9. Functions and Operators .................................................. 198
  9.1. Logical Operators ...................................................... 198
  9.2. Comparison Functions and Operators ............................... 198
  9.3. Mathematical Functions and Operators ............................. 201
  9.4. String Functions and Operators .................................... 204
    9.4.1. format .......................................................... 217
  9.5. Binary String Functions and Operators ............................ 219
  9.6. Bit String Functions and Operators ............................... 221
  9.7. Pattern Matching ................................................... 222
    9.7.1. LIKE .......................................................... 223
    9.7.2. SIMILAR TO Regular Expressions .............................. 223
    9.7.3. POSIX Regular Expressions .................................. 224
  9.8. Data Type Formatting Functions ................................... 237
  9.9. Date/Time Functions and Operators ................................ 244
    9.9.1. EXTRACT, date_part .......................................... 249
    9.9.2. date_trunc .................................................. 253
    9.9.3. AT TIME ZONE ............................................... 254
    9.9.4. Current Date/Time .......................................... 255
    9.9.5. Delaying Execution .......................................... 256
  9.10. Enum Support Functions ............................................ 257
  9.11. Geometric Functions and Operators ............................... 258
  9.12. Network Address Functions and Operators ......................... 262
  9.13. Text Search Functions and Operators ............................. 264
  9.14. XML Functions ..................................................... 271
    9.14.2. XML Predicates .............................................. 275
    9.14.3. Processing XML .............................................. 277
    9.14.4. Mapping Tables to XML ..................................... 281
  9.15. JSON Functions and Operators .................................... 284
  9.16. Sequence Manipulation Functions ................................ 293
  9.17. Conditional Expressions .......................................... 295
    9.17.1. CASE ..................................................... 296
    9.17.2. COALESCE .................................................. 297
    9.17.3. NULLIF ................................................... 297
    9.17.4. GREATEST and LEAST ....................................... 298
  9.18. Array Functions and Operators .................................... 298
  9.19. Range Functions and Operators .................................... 301
  9.20. Aggregate Functions .............................................. 303
  9.21. Window Functions .................................................. 310
  9.22. Subquery Expressions ............................................. 312
    9.22.1. EXISTS ................................................... 312
    9.22.2. IN ........................................................ 313
    9.22.3. NOT IN .................................................... 313
    9.22.4. ANY/SOME .................................................. 314
    9.22.5. ALL ....................................................... 314
    9.22.6. Single-row Comparison ..................................... 315
  9.23. Row and Array Comparisons ....................................... 315
    9.23.1. IN ........................................................ 315
    9.23.2. NOT IN .................................................... 316
    9.23.3. ANY/SOME (array) ......................................... 316
    9.23.4. ALL (array) ................................................ 316
    9.23.5. Row Constructor Comparison ................................ 317
    9.23.6. Composite Type Comparison .................................. 317
  9.24. Set Returning Functions .......................................... 318
  9.25. System Information Functions ...................................... 321
  9.26. System Administration Functions .................................. 338
15.4.1. Parallel Labeling for Functions and Aggregates ........................................ 466
Chapter 4. SQL Syntax

This chapter describes the syntax of SQL. It forms the foundation for understanding the following chapters which will go into detail about how SQL commands are applied to define and modify data.

We also advise users who are already familiar with SQL to read this chapter carefully because it contains several rules and concepts that are implemented inconsistently among SQL databases or that are specific to PostgreSQL.

4.1. Lexical Structure

SQL input consists of a sequence of commands. A command is composed of a sequence of tokens, terminated by a semicolon (";"). The end of the input stream also terminates a command. Which tokens are valid depends on the syntax of the particular command.

A token can be a key word, an identifier, a quoted identifier, a literal (or constant), or a special character symbol. Tokens are normally separated by whitespace (space, tab, newline), but need not be if there is no ambiguity (which is generally only the case if a special character is adjacent to some other token type).

For example, the following is (syntactically) valid SQL input:

```
SELECT * FROM MY_TABLE;
UPDATE MY_TABLE SET A = 5;
INSERT INTO MY_TABLE VALUES (3, 'hi there');
```

This is a sequence of three commands, one per line (although this is not required; more than one command can be on a line, and commands can usefully be split across lines).

Additionally, comments can occur in SQL input. They are not tokens, they are effectively equivalent to whitespace.

The SQL syntax is not very consistent regarding what tokens identify commands and which are operands or parameters. The first few tokens are generally the command name, so in the above example we would usually speak of a "SELECT", an "UPDATE", and an "INSERT" command. But for instance the UPDATE command always requires a SET token to appear in a certain position, and this particular variation of INSERT also requires a VALUES in order to be complete. The precise syntax rules for each command are described in Part VI.

4.1.1. Identifiers and Key Words

Tokens such as SELECT, UPDATE, or VALUES in the example above are examples of key words, that is, words that have a fixed meaning in the SQL language. The tokens MY_TABLE and A are examples of identifiers. They identify names of tables, columns, or other database objects, depending on the command they are used in. Therefore they are sometimes simply called “names”. Key words and identifiers have the same lexical structure, meaning that one cannot know whether a token is an identifier or a key word without knowing the language. A complete list of key words can be found in Appendix C.

SQL identifiers and key words must begin with a letter (a-z, but also letters with diacritical marks and non-Latin letters) or an underscore (_). Subsequent characters in an identifier or key word can be letters, underscores, digits (0-9), or dollar signs ($). Note that dollar signs are not allowed in identifiers according to the letter of the SQL standard, so their use might render applications less portable. The SQL standard will not define a key word that contains digits or starts or ends with an underscore, so identifiers of this form are safe against possible conflict with future extensions of the standard.
The system uses no more than `NAMEDATALEN-1` bytes of an identifier; longer names can be written in commands, but they will be truncated. By default, `NAMEDATALEN` is 64 so the maximum identifier length is 63 bytes. If this limit is problematic, it can be raised by changing the `NAMEDATALEN` constant in `src/include/pg_config_manual.h`.

Key words and unquoted identifiers are case insensitive. Therefore:

```
UPDATE MY_TABLE SET A = 5;
```

can equivalently be written as:

```
UPDTE my_TaBE SeT a = 5;
```

A convention often used is to write key words in upper case and names in lower case, e.g.:

```
UPDATE my_table SET a = 5;
```

There is a second kind of identifier: the delimited identifier or quoted identifier. It is formed by enclosing an arbitrary sequence of characters in double-quotes (`"`). A delimited identifier is always an identifier, never a key word. So "select" could be used to refer to a column or table named "select", whereas an unquoted `select` would be taken as a key word and would therefore provoke a parse error when used where a table or column name is expected. The example can be written with quoted identifiers like this:

```
UPDATE "my_table" SET "a" = 5;
```

Quoted identifiers can contain any character, except the character with code zero. (To include a double quote, write two double quotes.) This allows constructing table or column names that would otherwise not be possible, such as ones containing spaces or ampersands. The length limitation still applies.

A variant of quoted identifiers allows including escaped Unicode characters identified by their code points. This variant starts with `U&` (upper or lower case U followed by ampersand) immediately before the opening double quote, without any spaces in between, for example `U&"foo"`. (Note that this creates an ambiguity with the operator `&`. Use spaces around the operator to avoid this problem.) Inside the quotes, Unicode characters can be specified in escaped form by writing a backslash followed by the four-digit hexadecimal code point number or alternatively a backslash followed by a plus sign followed by a six-digit hexadecimal code point number. For example, the identifier "data" could be written as

```
U&"d\0061t\+000061"
```

The following less trivial example writes the Russian word “slon” (elephant) in Cyrillic letters:

```
U&"\0441\043B\043E\043D"
```

If a different escape character than backslash is desired, it can be specified using the `UESCAPE` clause after the string, for example:

```
U&"d!0061t!+000061" UESCAPE '!' 
```

The escape character can be any single character other than a hexadecimal digit, the plus sign, a single quote, a double quote, or a whitespace character. Note that the escape character is written in single quotes, not double quotes.
To include the escape character in the identifier literally, write it twice.

The Unicode escape syntax works only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \007F) can be specified. Both the 4-digit and the 6-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 6-digit form technically makes this unnecessary. (Surrogate pairs are not stored directly, but combined into a single code point that is then encoded in UTF-8.)

Quoting an identifier also makes it case-sensitive, whereas unquoted names are always folded to lower case. For example, the identifiers FOO, foo, and "foo" are considered the same by PostgreSQL, but "Foo" and "FOO" are different from these three and each other. (The folding of unquoted names to lower case in PostgreSQL is incompatible with the SQL standard, which says that unquoted names should be folded to upper case. Thus, foo should be equivalent to "FOO" not "foo" according to the standard. If you want to write portable applications you are advised to always quote a particular name or never quote it.)

4.1.2. Constants

There are three kinds of implicitly-typed constants in PostgreSQL: strings, bit strings, and numbers. Constants can also be specified with explicit types, which can enable more accurate representation and more efficient handling by the system. These alternatives are discussed in the following subsections.

4.1.2.1. String Constants

A string constant in SQL is an arbitrary sequence of characters bounded by single quotes ('), for example 'This is a string'. To include a single-quote character within a string constant, write two adjacent single quotes, e.g., 'Dianne''s horse'. Note that this is not the same as a double-quote character (").

Two string constants that are only separated by whitespace with at least one newline are concatenated and effectively treated as if the string had been written as one constant. For example:

```
SELECT 'foo'
     'bar';
```

is equivalent to:

```
SELECT 'foobar';
```

but:

```
SELECT 'foo'     'bar';
```

is not valid syntax. (This slightly bizarre behavior is specified by SQL; PostgreSQL is following the standard.)

4.1.2.2. String Constants with C-style Escapes

PostgreSQL also accepts “escape” string constants, which are an extension to the SQL standard. An escape string constant is specified by writing the letter E (upper or lower case) just before the opening single quote, e.g., E'foo'. (When continuing an escape string constant across lines, write E only before the first opening quote.) Within an escape string, a backslash character (\) begins a C-like backslash escape sequence, in which the combination of backslash and following character(s) represent a special byte value, as shown in Table 4.1.
Table 4.1. Backslash Escape Sequences

<table>
<thead>
<tr>
<th>Backslash Escape Sequence</th>
<th>Interpretation</th>
</tr>
</thead>
<tbody>
<tr>
<td>\b</td>
<td>backspace</td>
</tr>
<tr>
<td>\f</td>
<td>form feed</td>
</tr>
<tr>
<td>\n</td>
<td>newline</td>
</tr>
<tr>
<td>\r</td>
<td>carriage return</td>
</tr>
<tr>
<td>\t</td>
<td>tab</td>
</tr>
<tr>
<td>\o, \oo, \ooo (o = 0 - 7)</td>
<td>octal byte value</td>
</tr>
<tr>
<td>\xh, \xhh (h = 0 - 9, A - F)</td>
<td>hexadecimal byte value</td>
</tr>
<tr>
<td>\uxxxx, \Uxxxxxxxxx (x = 0 - 9, A - F)</td>
<td>16 or 32-bit hexadecimal Unicode character value</td>
</tr>
</tbody>
</table>

Any other character following a backslash is taken literally. Thus, to include a backslash character, write two backslashes (\\). Also, a single quote can be included in an escape string by writing \', in addition to the normal way of ’’.

It is your responsibility that the byte sequences you create, especially when using the octal or hexadecimal escapes, compose valid characters in the server character set encoding. When the server encoding is UTF-8, then the Unicode escapes or the alternative Unicode escape syntax, explained in Section 4.1.2.3, should be used instead. (The alternative would be doing the UTF-8 encoding by hand and writing out the bytes, which would be very cumbersome.)

The Unicode escape syntax works fully only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \u007F) can be specified. Both the 4-digit and the 8-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 8-digit form technically makes this unnecessary. (When surrogate pairs are used when the server encoding is UTF8, they are first combined into a single code point that is then encoded in UTF-8.)

Caution

If the configuration parameter standard_conforming_strings is off, then PostgreSQL recognizes backslash escapes in both regular and escape string constants. However, as of PostgreSQL 9.1, the default is on, meaning that backslash escapes are recognized only in escape string constants. This behavior is more standards-compliant, but might break applications which rely on the historical behavior, where backslash escapes were always recognized. As a workaround, you can set this parameter to off, but it is better to migrate away from using backslash escapes. If you need to use a backslash escape to represent a special character, write the string constant with an E.

In addition to standard_conforming_strings, the configuration parameters escape_string_warning and backslash_quote govern treatment of backslashes in string constants.

The character with the code zero cannot be in a string constant.

4.1.2.3. String Constants with Unicode Escapes

PostgreSQL also supports another type of escape syntax for strings that allows specifying arbitrary Unicode characters by code point. A Unicode escape string constant starts with \u (upper or lower case letter U followed by ampersand) immediately before the opening quote, without any spaces in between, for example \u'\u007F'. (Note that this creates an ambiguity with the operator \&. Use spaces around the operator to avoid this problem.) Inside the quotes, Unicode characters can be specified
in escaped form by writing a backslash followed by the four-digit hexadecimal code point number or alternatively a backslash followed by a plus sign followed by a six-digit hexadecimal code point number. For example, the string 'data' could be written as

U\'d\0061t\+000061'

The following less trivial example writes the Russian word “slon” (elephant) in Cyrillic letters:

U\'\0441\043B\043E\043D'

If a different escape character than backslash is desired, it can be specified using the UESCAPE clause after the string, for example:

U\'d!0061t!+000061' UESCAPE '!

The escape character can be any single character other than a hexadecimal digit, the plus sign, a single quote, a double quote, or a whitespace character.

The Unicode escape syntax works only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \007F) can be specified. Both the 4-digit and the 6-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 6-digit form technically makes this unnecessary. (When surrogate pairs are used when the server encoding is UTF8, they are first combined into a single code point that is then encoded in UTF-8.)

Also, the Unicode escape syntax for string constants only works when the configuration parameter standard_conforming_strings is turned on. This is because otherwise this syntax could confuse clients that parse the SQL statements to the point that it could lead to SQL injections and similar security issues. If the parameter is set to off, this syntax will be rejected with an error message.

To include the escape character in the string literally, write it twice.

### 4.1.2.4. Dollar-quoted String Constants

While the standard syntax for specifying string constants is usually convenient, it can be difficult to understand when the desired string contains many single quotes or backslashes, since each of those must be doubled. To allow more readable queries in such situations, PostgreSQL provides another way, called “dollar quoting”, to write string constants. A dollar-quoted string constant consists of a dollar sign ($), an optional “tag” of zero or more characters, another dollar sign, an arbitrary sequence of characters that makes up the string content, a dollar sign, the same tag that began this dollar quote, and a dollar sign. For example, here are two different ways to specify the string “Dianne's horse” using dollar quoting:

$$Dianne's horse$$
$SomeTag$Dianne's horse$SomeTag$

Notice that inside the dollar-quoted string, single quotes can be used without needing to be escaped. Indeed, no characters inside a dollar-quoted string are ever escaped: the string content is always written literally. Backslashes are not special, and neither are dollar signs, unless they are part of a sequence matching the opening tag.

It is possible to nest dollar-quoted string constants by choosing different tags at each nesting level. This is most commonly used in writing function definitions. For example:

$function$
BEGIN
    RETURN ($1 ~ $q$[\t\r\n\v\"]$q$);
END;
$function$

Here, the sequence $q$[\t\r\n\v\"]$q$ represents a dollar-quoted literal string \[\t\r\n\v\"\], which will be recognized when the function body is executed by PostgreSQL. But since the sequence does not match the outer dollar quoting delimiter $function$, it is just some more characters within the constant so far as the outer string is concerned.

The tag, if any, of a dollar-quoted string follows the same rules as an unquoted identifier, except that it cannot contain a dollar sign. Tags are case sensitive, so $tag$string content$tag$ is correct, but $TAG$string content$tag$ is not.

A dollar-quoted string that follows a keyword or identifier must be separated from it by whitespace; otherwise the dollar quoting delimiter would be taken as part of the preceding identifier.

Dollar quoting is not part of the SQL standard, but it is often a more convenient way to write complicated string literals than the standard-compliant single quote syntax. It is particularly useful when representing string constants inside other constants, as is often needed in procedural function definitions. With single-quote syntax, each backslash in the above example would have to be written as four backslashes, which would be reduced to two backslashes in parsing the original string constant, and then to one when the inner string constant is re-parsed during function execution.

### 4.1.2.5. Bit-string Constants

Bit-string constants look like regular string constants with a B (upper or lower case) immediately before the opening quote (no intervening whitespace), e.g., B'1001'. The only characters allowed within bit-string constants are 0 and 1.

Alternatively, bit-string constants can be specified in hexadecimal notation, using a leading X (upper or lower case), e.g., X'1FF'. This notation is equivalent to a bit-string constant with four binary digits for each hexadecimal digit.

Both forms of bit-string constant can be continued across lines in the same way as regular string constants. Dollar quoting cannot be used in a bit-string constant.

### 4.1.2.6. Numeric Constants

Numeric constants are accepted in these general forms:

- `digits`
- `digits.[digits][e[+-]digits]`
- `[digits].digits[e[+-]digits]`
- `digits[+-]digits`

where `digits` is one or more decimal digits (0 through 9). At least one digit must be before or after the decimal point, if one is used. At least one digit must follow the exponent marker (e), if one is present. There cannot be any spaces or other characters embedded in the constant. Note that any leading plus or minus sign is not actually considered part of the constant; it is an operator applied to the constant.

These are some examples of valid numeric constants:

- 42
- 3.5
- 4.
- .001
A numeric constant that contains neither a decimal point nor an exponent is initially presumed to be type integer if its value fits in type integer (32 bits); otherwise it is presumed to be type bigint if its value fits in type bigint (64 bits); otherwise it is taken to be type numeric. Constants that contain decimal points and/or exponents are always initially presumed to be type numeric.

The initially assigned data type of a numeric constant is just a starting point for the type resolution algorithms. In most cases the constant will be automatically coerced to the most appropriate type depending on context. When necessary, you can force a numeric value to be interpreted as a specific data type by casting it. For example, you can force a numeric value to be treated as type real(float4) by writing:

```sql
REAL '1.23'  -- string style
1.23::REAL   -- PostgreSQL (historical) style
```

These are actually just special cases of the general casting notations discussed next.

### 4.1.2.7. Constants of Other Types

A constant of an arbitrary type can be entered using any one of the following notations:

- `type 'string'`
- `'string'::type`
- `CAST ( 'string' AS type )`

The string constant's text is passed to the input conversion routine for the type called `type`. The result is a constant of the indicated type. The explicit type cast can be omitted if there is no ambiguity as to the type the constant must be (for example, when it is assigned directly to a table column), in which case it is automatically coerced.

The string constant can be written using either regular SQL notation or dollar-quoting.

It is also possible to specify a type coercion using a function-like syntax:

```sql
typename ( 'string' )
```

but not all type names can be used in this way; see Section 4.2.9 for details.

The `::`, `CAST()`, and function-call syntaxes can also be used to specify run-time type conversions of arbitrary expressions, as discussed in Section 4.2.9. To avoid syntactic ambiguity, the `type 'string'` syntax can only be used to specify the type of a simple literal constant. Another restriction on the `type 'string'` syntax is that it does not work for array types; use `::` or `CAST()` to specify the type of an array constant.

The `CAST()` syntax conforms to SQL. The `type 'string'` syntax is a generalization of the standard: SQL specifies this syntax only for a few data types, but PostgreSQL allows it for all types. The syntax with `::` is historical PostgreSQL usage, as is the function-call syntax.

### 4.1.3. Operators

An operator name is a sequence of up to NAMEDATALEN-1 (63 by default) characters from the following list:
+ - * / < > = ~ ! @ # % ^ & | ` ?

There are a few restrictions on operator names, however:

- `--` and `/` cannot appear anywhere in an operator name, since they will be taken as the start of a comment.
- A multiple-character operator name cannot end in `+` or `–`, unless the name also contains at least one of these characters: ~ ! @ # % ^ & | ` ?

For example, `@–` is an allowed operator name, but `*–` is not. This restriction allows PostgreSQL to parse SQL-compliant queries without requiring spaces between tokens.

When working with non-SQL-standard operator names, you will usually need to separate adjacent operators with spaces to avoid ambiguity. For example, if you have defined a left unary operator named `@`, you cannot write `X*@Y`; you must write `X* @Y` to ensure that PostgreSQL reads it as two operator names not one.

### 4.1.4. Special Characters

Some characters that are not alphanumeric have a special meaning that is different from being an operator. Details on the usage can be found at the location where the respective syntax element is described. This section only exists to advise the existence and summarize the purposes of these characters.

- A dollar sign (`$`) followed by digits is used to represent a positional parameter in the body of a function definition or a prepared statement. In other contexts the dollar sign can be part of an identifier or a dollar-quoted string constant.
- Parentheses (`()`) have their usual meaning to group expressions and enforce precedence. In some cases parentheses are required as part of the fixed syntax of a particular SQL command.
- Brackets (`[ ]`) are used to select the elements of an array. See Section 8.15 for more information on arrays.
- Commas (`,`) are used in some syntactical constructs to separate the elements of a list.
- The semicolon (`;`) terminates an SQL command. It cannot appear anywhere within a command, except within a string constant or quoted identifier.
- The colon (`:`) is used to select “slices” from arrays. (See Section 8.15.) In certain SQL dialects (such as Embedded SQL), the colon is used to prefix variable names.
- The asterisk (`*`) is used in some contexts to denote all the fields of a table row or composite value. It also has a special meaning when used as the argument of an aggregate function, namely that the aggregate does not require any explicit parameter.
- The period (`.`) is used in numeric constants, and to separate schema, table, and column names.

### 4.1.5. Comments

A comment is a sequence of characters beginning with double dashes and extending to the end of the line, e.g.:

```
-- This is a standard SQL comment
```

Alternatively, C-style block comments can be used:
/* multiline comment
   * with nesting: /* nested block comment */
*/

where the comment begins with /* and extends to the matching occurrence of */. These block comments nest, as specified in the SQL standard but unlike C, so that one can comment out larger blocks of code that might contain existing block comments.

A comment is removed from the input stream before further syntax analysis and is effectively replaced by whitespace.

### 4.1.6. Operator Precedence

Table 4.2 shows the precedence and associativity of the operators in PostgreSQL. Most operators have the same precedence and are left-associative. The precedence and associativity of the operators is hard-wired into the parser.

You will sometimes need to add parentheses when using combinations of binary and unary operators. For instance:

```sql
SELECT 5 ! - 6;
```

will be parsed as:

```sql
SELECT 5 ! (- 6);
```

because the parser has no idea — until it is too late — that ! is defined as a postfix operator, not an infix one. To get the desired behavior in this case, you must write:

```sql
SELECT (5 !) - 6;
```

This is the price one pays for extensibility.

<table>
<thead>
<tr>
<th>Operator/Element</th>
<th>Associativity</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>.</td>
<td>left</td>
<td>table/column name separator</td>
</tr>
<tr>
<td>::</td>
<td>left</td>
<td>PostgreSQL-style typecast</td>
</tr>
<tr>
<td>[ ]</td>
<td>left</td>
<td>array element selection</td>
</tr>
<tr>
<td>+ -</td>
<td>right</td>
<td>unary plus, unary minus</td>
</tr>
<tr>
<td>^</td>
<td>left</td>
<td>exponentiation</td>
</tr>
<tr>
<td>* / %</td>
<td>left</td>
<td>multiplication, division, modulo</td>
</tr>
<tr>
<td>+ -</td>
<td>left</td>
<td>addition, subtraction</td>
</tr>
<tr>
<td>(any other operator)</td>
<td>left</td>
<td>all other native and user-defined operators</td>
</tr>
<tr>
<td>BETWEEN IN LIKE ILIKE SIMILAR</td>
<td>left</td>
<td>range containment, set membership, string matching</td>
</tr>
<tr>
<td>&lt;&gt; &lt;= =&gt; &gt;= &lt;&gt;</td>
<td></td>
<td>comparison operators</td>
</tr>
<tr>
<td>IS ISNULL NOTNULL</td>
<td></td>
<td>IS TRUE, IS FALSE, IS NULL, IS DISTINCT FROM, etc</td>
</tr>
<tr>
<td>Operator/Element</td>
<td>Associativity</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>---------------</td>
<td>--------------------</td>
</tr>
<tr>
<td>NOT</td>
<td>right</td>
<td>logical negation</td>
</tr>
<tr>
<td>AND</td>
<td>left</td>
<td>logical conjunction</td>
</tr>
<tr>
<td>OR</td>
<td>left</td>
<td>logical disjunction</td>
</tr>
</tbody>
</table>

Note that the operator precedence rules also apply to user-defined operators that have the same names as the built-in operators mentioned above. For example, if you define a “+” operator for some custom data type it will have the same precedence as the built-in “+” operator, no matter what yours does.

When a schema-qualified operator name is used in the `OPERATOR` syntax, as for example in:

```
SELECT 3 OPERATOR(pg_catalog.+) 4;
```

the `OPERATOR` construct is taken to have the default precedence shown in Table 4.2 for “any other operator”. This is true no matter which specific operator appears inside `OPERATOR()`.

**Note**

PostgreSQL versions before 9.5 used slightly different operator precedence rules. In particular, `<=` and `<>` used to be treated as generic operators; `IS` tests used to have higher priority; and `NOT BETWEEN` and related constructs acted inconsistently, being taken in some cases as having the precedence of `NOT` rather than `BETWEEN`. These rules were changed for better compliance with the SQL standard and to reduce confusion from inconsistent treatment of logically equivalent constructs. In most cases, these changes will result in no behavioral change, or perhaps in “no such operator” failures which can be resolved by adding parentheses. However, there are corner cases in which a query might change behavior without any parsing error being reported. If you are concerned about whether these changes have silently broken something, you can test your application with the configuration parameter `operator_precedence_warning` turned on to see if any warnings are logged.

### 4.2. Value Expressions

Value expressions are used in a variety of contexts, such as in the target list of the `SELECT` command, as new column values in `INSERT` or `UPDATE`, or in search conditions in a number of commands. The result of a value expression is sometimes called a **scalar**, to distinguish it from the result of a table expression (which is a table). Value expressions are therefore also called **scalar expressions** (or even simply **expressions**). The expression syntax allows the calculation of values from primitive parts using arithmetic, logical, set, and other operations.

A value expression is one of the following:

- A constant or literal value
- A column reference
- A positional parameter reference, in the body of a function definition or prepared statement
- A subscripted expression
- A field selection expression
- An operator invocation
- A function call
• An aggregate expression
• A window function call
• A type cast
• A collation expression
• A scalar subquery
• An array constructor
• A row constructor
• Another value expression in parentheses (used to group subexpressions and override precedence)

In addition to this list, there are a number of constructs that can be classified as an expression but do not follow any general syntax rules. These generally have the semantics of a function or operator and are explained in the appropriate location in Chapter 9. An example is the IS NULL clause.

We have already discussed constants in Section 4.1.2. The following sections discuss the remaining options.

4.2.1. Column References

A column can be referenced in the form:

correlation.columnname

correlation is the name of a table (possibly qualified with a schema name), or an alias for a table defined by means of a FROM clause. The correlation name and separating dot can be omitted if the column name is unique across all the tables being used in the current query. (See also Chapter 7.)

4.2.2. Positional Parameters

A positional parameter reference is used to indicate a value that is supplied externally to an SQL statement. Parameters are used in SQL function definitions and in prepared queries. Some client libraries also support specifying data values separately from the SQL command string, in which case parameters are used to refer to the out-of-line data values. The form of a parameter reference is:

$number

For example, consider the definition of a function, dept, as:

CREATE FUNCTION dept(text) RETURNS dept
  AS $$ SELECT * FROM dept WHERE name = $1 $$
  LANGUAGE SQL;

Here the $1 references the value of the first function argument whenever the function is invoked.

4.2.3. Subscripts

If an expression yields a value of an array type, then a specific element of the array value can be extracted by writing

expression[subscript]
or multiple adjacent elements (an “array slice”) can be extracted by writing

\[expression[lower_subscript:upper_subscript]\]

(Here, the brackets \[ \] are meant to appear literally.) Each \textit{subscript} is itself an expression, which must yield an integer value.

In general the array \textit{expression} must be parenthesized, but the parentheses can be omitted when the expression to be subscripted is just a column reference or positional parameter. Also, multiple subscripts can be concatenated when the original array is multidimensional. For example:

\[
\begin{align*}
\text{mytable.arraycolumn[4]} \\
\text{mytable.two_d_column[17][34]} \\
\text{$1[10:42]} \\
\text{(arrayfunction(a,b))[42]}
\end{align*}
\]

The parentheses in the last example are required. See Section 8.15 for more about arrays.

### 4.2.4. Field Selection

If an expression yields a value of a composite type (row type), then a specific field of the row can be extracted by writing

\[expression.fieldname\]

In general the row \textit{expression} must be parenthesized, but the parentheses can be omitted when the expression to be selected from is just a table reference or positional parameter. For example:

\[
\begin{align*}
\text{mytable.mycolumn} \\
\text{$1.somecolumn} \\
\text{(rowfunction(a,b)).col3}
\end{align*}
\]

(Thus, a qualified column reference is actually just a special case of the field selection syntax.) An important special case is extracting a field from a table column that is of a composite type:

\[
\begin{align*}
\text{(compositecol).somefield} \\
\text{(mytable.compositecol).somefield}
\end{align*}
\]

The parentheses are required here to show that \textit{compositecol} is a column name not a table name, or that \textit{mytable} is a table name not a schema name in the second case.

You can ask for all fields of a composite value by writing \texttt{.*}:

\[
\begin{align*}
\text{(compositecol).*}
\end{align*}
\]

This notation behaves differently depending on context; see Section 8.16.5 for details.

### 4.2.5. Operator Invocations

There are three possible syntaxes for an operator invocation:

\[
\begin{align*}
\text{expression operator expression (binary infix operator)} \\
\text{operator expression (unary prefix operator)}
\end{align*}
\]
expression operator (unary postfix operator)

where the operator token follows the syntax rules of Section 4.1.3, or is one of the key words AND, OR, and NOT, or is a qualified operator name in the form:

OPERATOR(schema_operatorname)

Which particular operators exist and whether they are unary or binary depends on what operators have been defined by the system or the user. Chapter 9 describes the built-in operators.

4.2.6. Function Calls

The syntax for a function call is the name of a function (possibly qualified with a schema name), followed by its argument list enclosed in parentheses:

function_name ([expression [, expression ... ]])

For example, the following computes the square root of 2:

sqrt(2)

The list of built-in functions is in Chapter 9. Other functions can be added by the user.

When issuing queries in a database where some users mistrust other users, observe security precautions from Section 10.3 when writing function calls.

The arguments can optionally have names attached. See Section 4.3 for details.

Note

A function that takes a single argument of composite type can optionally be called using field-selection syntax, and conversely field selection can be written in functional style. That is, the notations col(table) and table.col are interchangeable. This behavior is not SQL-standard but is provided in PostgreSQL because it allows use of functions to emulate “computed fields”. For more information see Section 8.16.5.

4.2.7. Aggregate Expressions

An aggregate expression represents the application of an aggregate function across the rows selected by a query. An aggregate function reduces multiple inputs to a single output value, such as the sum or average of the inputs. The syntax of an aggregate expression is one of the following:

\[
\text{aggregate_name} (\text{expression} [ , \ldots ] \ [ \text{order_by_clause} ] \ [ \text{FILTER} \ (\text{WHERE} \ \text{filter_clause})]) \\
\text{aggregate_name} (\text{ALL} \ \text{expression} [ , \ldots ] \ [ \text{order_by_clause} ] \ [ \text{FILTER} \ (\text{WHERE} \ \text{filter_clause})]) \\
\text{aggregate_name} (\text{DISTINCT} \ \text{expression} [ , \ldots ] \ [ \text{order_by_clause} ] \ [ \text{FILTER} \ (\text{WHERE} \ \text{filter_clause})]) \\
\text{aggregate_name} (\ast) \ [ \text{FILTER} \ (\text{WHERE} \ \text{filter_clause})] \\
\text{aggregate_name} ([ \text{expression} [, \ldots ] ]) \ \text{WITHIN} \ \text{GROUP} \ (\text{order_by_clause}) \ [ \text{FILTER} \ (\text{WHERE} \ \text{filter_clause})]
\]

where aggregate_name is a previously defined aggregate (possibly qualified with a schema name) and expression is any value expression that does not itself contain an aggregate expression or
a window function call. The optional `order_by_clause` and `filter_clause` are described below.

The first form of aggregate expression invokes the aggregate once for each input row. The second form is the same as the first, since `ALL` is the default. The third form invokes the aggregate once for each distinct value of the expression (or distinct set of values, for multiple expressions) found in the input rows. The fourth form invokes the aggregate once for each input row; since no particular input value is specified, it is generally only useful for the `count(*)` aggregate function. The last form is used with `ordered-set` aggregate functions, which are described below.

Most aggregate functions ignore null inputs, so that rows in which one or more of the expression(s) yield null are discarded. This can be assumed to be true, unless otherwise specified, for all built-in aggregates.

For example, `count(*)` yields the total number of input rows; `count(f1)` yields the number of input rows in which `f1` is non-null, since `count` ignores nulls; and `count(distinct f1)` yields the number of distinct non-null values of `f1`.

Ordinarily, the input rows are fed to the aggregate function in an unspecified order. In many cases this does not matter; for example, `min` produces the same result no matter what order it receives the inputs in. However, some aggregate functions (such as `array_agg` and `string_agg`) produce results that depend on the ordering of the input rows. When using such an aggregate, the optional `order_by_clause` can be used to specify the desired ordering. The `order_by_clause` has the same syntax as for a query-level `ORDER BY` clause, as described in Section 7.5, except that its expressions are always just expressions and cannot be output-column names or numbers. For example:

```sql
SELECT array_agg(a ORDER BY b DESC) FROM table;
```

When dealing with multiple-argument aggregate functions, note that the `ORDER BY` clause goes after all the aggregate arguments. For example, write this:

```sql
SELECT string_agg(a, ',') ORDER BY a) FROM table;
```

not this:

```sql
SELECT string_agg(a ORDER BY a, ',') FROM table;  -- incorrect
```

The latter is syntactically valid, but it represents a call of a single-argument aggregate function with two `ORDER BY` keys (the second one being rather useless since it's a constant).

If `DISTINCT` is specified in addition to an `order_by_clause`, then all the `ORDER BY` expressions must match regular arguments of the aggregate; that is, you cannot sort on an expression that is not included in the `DISTINCT` list.

---

**Note**

The ability to specify both `DISTINCT` and `ORDER BY` in an aggregate function is a PostgreSQL extension.

Placing `ORDER BY` within the aggregate's regular argument list, as described so far, is used when ordering the input rows for general-purpose and statistical aggregates, for which ordering is optional. There is a subclass of aggregate functions called `ordered-set aggregates` for which an `order_by_clause` is required, usually because the aggregate's computation is only sensible in terms of a specific ordering of its input rows. Typical examples of ordered-set aggregates include rank and percentile calculations. For an ordered-set aggregate, the `order_by_clause` is written inside
WTIEN_GROUP (...), as shown in the final syntax alternative above. The expressions in the
order_by_clause are evaluated once per input row just like regular aggregate arguments, sorted
as per the order_by_clause's requirements, and fed to the aggregate function as input arguments.
(This is unlike the case for a non-WITHIN_GROUP order_by_clause, which is not treated as
argument(s) to the aggregate function.) The argument expressions preceding WITHIN_GROUP, if
any, are called direct arguments to distinguish them from the aggregated arguments listed in the or-
der_by_clause. Unlike regular aggregate arguments, direct arguments are evaluated only once
per aggregate call, not once per input row. This means that they can contain variables only if those
variables are grouped by GROUP BY; this restriction is the same as if the direct arguments were not
inside an aggregate expression at all. Direct arguments are typically used for things like percentile
fractions, which only make sense as a single value per aggregation calculation. The direct argument
list can be empty; in this case, write just () not (*). (PostgreSQL will actually accept either spelling,
but only the first way conforms to the SQL standard.)

An example of an ordered-set aggregate call is:

```
SELECT percentile_cont(0.5) WITHIN GROUP (ORDER BY income) FROM households;
percentile_cont
-----------
      50489
```

which obtains the 50th percentile, or median, value of the income column from table households.
Here, 0.5 is a direct argument; it would make no sense for the percentile fraction to be a value varying
across rows.

If FILTER is specified, then only the input rows for which the filter_clause evaluates to true
are fed to the aggregate function; other rows are discarded. For example:

```
SELECT count(*) AS unfiltered,
        count(*) FILTER (WHERE i < 5) AS filtered
FROM generate_series(1,10) AS s(i);
    unfiltered | filtered
-------------+----------
          10 |         4
(1 row)
```

The predefined aggregate functions are described in Section 9.20. Other aggregate functions can be
added by the user.

An aggregate expression can only appear in the result list or HAVING clause of a SELECT command.
It is forbidden in other clauses, such as WHERE, because those clauses are logically evaluated before
the results of aggregates are formed.

When an aggregate expression appears in a subquery (see Section 4.2.11 and Section 9.22), the aggre-
gate is normally evaluated over the rows of the subquery. But an exception occurs if the aggregate's
arguments (and filter_clause if any) contain only outer-level variables: the aggregate then be-
longs to the nearest such outer level, and is evaluated over the rows of that query. The aggregate ex-
pression as a whole is then an outer reference for the subquery it appears in, and acts as a constant over
any one evaluation of that subquery. The restriction about appearing only in the result list or HAVING
clause applies with respect to the query level that the aggregate belongs to.

**4.2.8. Window Function Calls**

A window function call represents the application of an aggregate-like function over some portion of
the rows selected by a query. Unlike non-window aggregate calls, this is not tied to grouping of the
selected rows into a single output row — each row remains separate in the query output. However the window function has access to all the rows that would be part of the current row’s group according to the grouping specification (PARTITION BY list) of the window function call. The syntax of a window function call is one of the following:

\[
\text{function_name} \left( \left[ \text{expression} \left[ , \text{expression} \ldots \right] \right] \right) \left[ \text{FILTER} \left( \text{WHERE filter_clause} \right) \right] \text{OVER} \; \text{window_name} \\
\text{function_name} \left( \left[ \text{expression} \left[ , \text{expression} \ldots \right] \right] \right) \left[ \text{FILTER} \left( \text{WHERE filter_clause} \right) \right] \text{OVER} \left( \text{window_definition} \right) \\
\text{function_name} \left( \ast \right) \left[ \text{FILTER} \left( \text{WHERE filter_clause} \right) \right] \text{OVER} \; \text{window_name} \\
\text{function_name} \left( \ast \right) \left[ \text{FILTER} \left( \text{WHERE filter_clause} \right) \right] \text{OVER} \left( \text{window_definition} \right)
\]

where \text{window_definition} has the syntax

\[
\left[ \text{existing_window_name} \right] \left[ \text{PARTITION BY expression} \left[ , \ldots \right] \right] \left[ \text{ORDER BY expression} \left[ \text{ASC} | \text{DESC} | \text{USING operator} \right] \left[ \text{NULLS} \left( \text{FIRST} | \text{LAST} \right) \right] \left[ , \ldots \right] \right] \left[ \text{frame_clause} \right]
\]

The optional \text{frame_clause} can be one of

\[
\{ \text{RANGE} | \text{ROWS} | \text{GROUPS} \} \text{frame_start} \left[ \text{frame_exclusion} \right] \\
\{ \text{RANGE} | \text{ROWS} | \text{GROUPS} \} \text{BETWEEN} \text{frame_start} \text{AND} \text{frame_end} \left[ \text{frame_exclusion} \right]
\]

where \text{frame_start} and \text{frame_end} can be one of

UNBOUNDED PRECEDING
offset PRECEDING
CURRENT ROW
offset FOLLOWING
UNBOUNDED FOLLOWING

and \text{frame_exclusion} can be one of

EXCLUDE CURRENT ROW
EXCLUDE GROUP
EXCLUDE TIES
EXCLUDE NO OTHERS

Here, \text{expression} represents any value expression that does not itself contain window function calls.

\text{window_name} is a reference to a named window specification defined in the query’s \text{WINDOW} clause. Alternatively, a full \text{window_definition} can be given within parentheses, using the same syntax as for defining a named window in the \text{WINDOW} clause; see the SELECT reference page for details. It’s worth pointing out that \text{OVER wname} is not exactly equivalent to \text{OVER (wname ...)}; the latter implies copying and modifying the window definition, and will be rejected if the referenced window specification includes a frame clause.

The PARTITION BY clause groups the rows of the query into \text{partitions}, which are processed separately by the window function. PARTITION BY works similarly to a query-level \text{GROUP BY} clause,
except that its expressions are always just expressions and cannot be output-column names or numbers. Without PARTITION BY, all rows produced by the query are treated as a single partition. The ORDER BY clause determines the order in which the rows of a partition are processed by the window function. It works similarly to a query-level ORDER BY clause, but likewise cannot use output-column names or numbers. Without ORDER BY, rows are processed in an unspecified order.

The frame_clause specifies the set of rows constituting the window frame, which is a subset of the current partition, for those window functions that act on the frame instead of the whole partition. The set of rows in the frame can vary depending on which row is the current row. The frame can be specified in RANGE, ROWS or GROUPS mode; in each case, it runs from the frame_start to the frame_end. If frame_end is omitted, the end defaults to CURRENT ROW.

A frame_start of UNBOUNDED PRECEDING means that the frame starts with the first row of the partition, and similarly a frame_end of UNBOUNDED FOLLOWING means that the frame ends with the last row of the partition.

In RANGE or GROUPS mode, a frame_start of CURRENT ROW means the frame starts with the current row’s first peer row (a row that the window’s ORDER BY clause sorts as equivalent to the current row), while a frame_end of CURRENT ROW means the frame ends with the current row’s last peer row. In ROWS mode, CURRENT ROW simply means the current row.

In the offset PRECEDING and offset FOLLOWING frame options, the offset must be an expression not containing any variables, aggregate functions, or window functions. The meaning of the offset depends on the frame mode:

- **In ROWS mode**, the offset must yield a non-null, non-negative integer, and the option means that the frame starts or ends the specified number of rows before or after the current row.

- **In GROUPS mode**, the offset again must yield a non-null, non-negative integer, and the option means that the frame starts or ends the specified number of peer groups before or after the current row’s peer group, where a peer group is a set of rows that are equivalent in the ORDER BY ordering. (There must be an ORDER BY clause in the window definition to use GROUPS mode.)

- **In RANGE mode**, these options require that the ORDER BY clause specify exactly one column. The offset specifies the maximum difference between the value of that column in the current row and its value in preceding or following rows of the frame. The data type of the offset expression varies depending on the data type of the ordering column. For numeric ordering columns it is typically of the same type as the ordering column, but for datetime ordering columns it is an interval. For example, if the ordering column is of type date or timestamp, one could write RANGE BETWEEN '1 day' PRECEDING AND '10 days' FOLLOWING. The offset is still required to be non-null and non-negative, though the meaning of “non-negative” depends on its data type.

In any case, the distance to the end of the frame is limited by the distance to the end of the partition, so that for rows near the partition ends the frame might contain fewer rows than elsewhere.

Notice that in both ROWS and GROUPS mode, 0 PRECEDING and 0 FOLLOWING are equivalent to CURRENT ROW. This normally holds in RANGE mode as well, for an appropriate data-type-specific meaning of “zero”.

The frame_exclusion option allows rows around the current row to be excluded from the frame, even if they would be included according to the frame start and frame end options. EXCLUDE CURRENT ROW excludes the current row from the frame. EXCLUDE GROUP excludes the current row and its ordering peers from the frame. EXCLUDE TIES excludes any peers of the current row from the frame, but not the current row itself. EXCLUDE NO OTHERS simply specifies explicitly the default behavior of not excluding the current row or its peers.

The default framing option is RANGE UNBOUNDED PRECEDING, which is the same as RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW. With ORDER BY, this sets the frame to be all rows from the partition start up through the current row’s last ORDER BY peer. Without
ORDER BY, this means all rows of the partition are included in the window frame, since all rows become peers of the current row.

Restrictions are that frame_start cannot be UNBOUNDED FOLLOWING, frame_end cannot be UNBOUNDED PRECEDING, and the frame_end choice cannot appear earlier in the above list of frame_start and frame_end options than the frame_start choice does — for example RANGE BETWEEN CURRENT ROW AND offset PRECEDING is not allowed. But, for example, ROWS BETWEEN 7 PRECEDING AND 8 PRECEDING is allowed, even though it would never select any rows.

If FILTER is specified, then only the input rows for which the filter_clause evaluates to true are fed to the window function; other rows are discarded. Only window functions that are aggregates accept a FILTER clause.

The built-in window functions are described in Table 9.57. Other window functions can be added by the user. Also, any built-in or user-defined general-purpose or statistical aggregate can be used as a window function. (Ordered-set and hypothetical-set aggregates cannot presently be used as window functions.)

The syntaxes using * are used for calling parameter-less aggregate functions as window functions, for example count(*) OVER (PARTITION BY x ORDER BY y). The asterisk (*) is customarily not used for window-specific functions. Window-specific functions do not allow DISTINCT or ORDER BY to be used within the function argument list.

Window function calls are permitted only in the SELECT list and the ORDER BY clause of the query.

More information about window functions can be found in Section 3.5, Section 9.21, and Section 7.2.5.

### 4.2.9. Type Casts

A type cast specifies a conversion from one data type to another. PostgreSQL accepts two equivalent syntaxes for type casts:

```sql
CAST ( expression AS type )
expression::type
```

The CAST syntax conforms to SQL; the syntax with :: is historical PostgreSQL usage.

When a cast is applied to a value expression of a known type, it represents a run-time type conversion. The cast will succeed only if a suitable type conversion operation has been defined. Notice that this is subtly different from the use of casts with constants, as shown in Section 4.1.2.7. A cast applied to an unadorned string literal represents the initial assignment of a type to a literal constant value, and so it will succeed for any type (if the contents of the string literal are acceptable input syntax for the data type).

An explicit type cast can usually be omitted if there is no ambiguity as to the type that a value expression must produce (for example, when it is assigned to a table column); the system will automatically apply a type cast in such cases. However, automatic casting is only done for casts that are marked “OK to apply implicitly” in the system catalogs. Other casts must be invoked with explicit casting syntax. This restriction is intended to prevent surprising conversions from being applied silently.

It is also possible to specify a type cast using a function-like syntax:

```sql
typename ( expression )
```

However, this only works for types whose names are also valid as function names. For example, double precision cannot be used this way, but the equivalent float8 can. Also, the names in-
interval, time, and timestamp can only be used in this fashion if they are double-quoted, because of syntactic conflicts. Therefore, the use of the function-like cast syntax leads to inconsistencies and should probably be avoided.

**Note**

The function-like syntax is in fact just a function call. When one of the two standard cast syntaxes is used to do a run-time conversion, it will internally invoke a registered function to perform the conversion. By convention, these conversion functions have the same name as their output type, and thus the "function-like syntax" is nothing more than a direct invocation of the underlying conversion function. Obviously, this is not something that a portable application should rely on. For further details see CREATE CAST.

### 4.2.10. Collation Expressions

The `COLLATE` clause overrides the collation of an expression. It is appended to the expression it applies to:

```
expr COLLATE collation
```

where `collation` is a possibly schema-qualified identifier. The `COLLATE` clause binds tighter than operators; parentheses can be used when necessary.

If no collation is explicitly specified, the database system either derives a collation from the columns involved in the expression, or it defaults to the default collation of the database if no column is involved in the expression.

The two common uses of the `COLLATE` clause are overriding the sort order in an `ORDER BY` clause, for example:

```
SELECT a, b, c FROM tbl WHERE ... ORDER BY a COLLATE "C";
```

and overriding the collation of a function or operator call that has locale-sensitive results, for example:

```
SELECT * FROM tbl WHERE a > 'foo' COLLATE "C";
```

Note that in the latter case the `COLLATE` clause is attached to an input argument of the operator we wish to affect. It doesn't matter which argument of the operator or function call the `COLLATE` clause is attached to, because the collation that is applied by the operator or function is derived by considering all arguments, and an explicit `COLLATE` clause will override the collations of all other arguments. (Attaching non-matching `COLLATE` clauses to more than one argument, however, is an error. For more details see Section 23.2.) Thus, this gives the same result as the previous example:

```
SELECT * FROM tbl WHERE a COLLATE "C" > 'foo';
```

But this is an error:

```
SELECT * FROM tbl WHERE (a > 'foo') COLLATE "C";
```

because it attempts to apply a collation to the result of the `>` operator, which is of the non-collatable data type `boolean`. 

50
4.2.11. Scalar Subqueries

A scalar subquery is an ordinary SELECT query in parentheses that returns exactly one row with one column. (See Chapter 7 for information about writing queries.) The SELECT query is executed and the single returned value is used in the surrounding value expression. It is an error to use a query that returns more than one row or more than one column as a scalar subquery. (But if, during a particular execution, the subquery returns no rows, there is no error; the scalar result is taken to be null.) The subquery can refer to variables from the surrounding query, which will act as constants during any one evaluation of the subquery. See also Section 9.22 for other expressions involving subqueries.

For example, the following finds the largest city population in each state:

```
SELECT name, (SELECT max(pop) FROM cities WHERE cities.state = states.name)
FROM states;
```

4.2.12. Array Constructors

An array constructor is an expression that builds an array value using values for its member elements. A simple array constructor consists of the keyword ARRAY, a left square bracket [, a list of expressions (separated by commas) for the array element values, and finally a right square bracket ]. For example:

```
SELECT ARRAY[1,2,3+4];
array
--------
{1,2,7}
(1 row)
```

By default, the array element type is the common type of the member expressions, determined using the same rules as for UNION or CASE constructs (see Section 10.5). You can override this by explicitly casting the array constructor to the desired type, for example:

```
SELECT ARRAY[1,2,22.7]::integer[];
array
---------
{1,2,23}
(1 row)
```

This has the same effect as casting each expression to the array element type individually. For more on casting, see Section 4.2.9.

Multidimensional array values can be built by nesting array constructors. In the inner constructors, the keyword ARRAY can be omitted. For example, these produce the same result:

```
SELECT ARRAY[ARRAY[1,2], ARRAY[3,4]];
array
---------------
{{1,2},{3,4}}
(1 row)
SELECT ARRAY[[1,2],[3,4]];
array
---------------
{{1,2},{3,4}}
(1 row)
```
Since multidimensional arrays must be rectangular, inner constructors at the same level must produce sub-arrays of identical dimensions. Any cast applied to the outer ARRAY constructor propagates automatically to all the inner constructors.

Multidimensional array constructor elements can be anything yielding an array of the proper kind, not only a sub-ARRAY construct. For example:

```
CREATE TABLE arr(f1 int[], f2 int[]);
INSERT INTO arr VALUES (ARRAY[[1,2],[3,4]], ARRAY[[5,6],[7,8]]);
SELECT ARRAY[f1, f2,('{{9,10},{11,12}})::int[]] FROM arr;
array
------------------------------------------------
{[1,2],[3,4]},{[5,6],[7,8]},{[9,10],[11,12]}
(1 row)
```

You can construct an empty array, but since it's impossible to have an array with no type, you must explicitly cast your empty array to the desired type. For example:

```
SELECT ARRAY[]::integer[];
array
-------
{}  
(1 row)
```

It is also possible to construct an array from the results of a subquery. In this form, the array constructor is written with the key word ARRAY followed by a parenthesized (not bracketed) subquery. For example:

```
SELECT ARRAY(SELECT oid FROM pg_proc WHERE proname LIKE 'bytea%');
array
-----------------------------------------------------------------------
(1 row)
SELECT ARRAY(SELECT ARRAY[i, i*2] FROM generate_series(1,5) AS a(i));
array
----------------------------------
{[1,2],[2,4],[3,6],[4,8],[5,10]}
(1 row)
```

The subquery must return a single column. If the subquery's output column is of a non-array type, the resulting one-dimensional array will have an element for each row in the subquery result, with an element type matching that of the subquery's output column. If the subquery's output column is of an array type, the result will be an array of the same type but one higher dimension; in this case all the subquery rows must yield arrays of identical dimensionality, else the result would not be rectangular.

The subscripts of an array value built with ARRAY always begin with one. For more information about arrays, see Section 8.15.

### 4.2.13. Row Constructors

A row constructor is an expression that builds a row value (also called a composite value) using values for its member fields. A row constructor consists of the key word ROW, a left parenthesis, zero or
more expressions (separated by commas) for the row field values, and finally a right parenthesis. For example:

```sql
SELECT ROW(1,2.5,'this is a test');
```

The key word `ROW` is optional when there is more than one expression in the list.

A row constructor can include the syntax `rowvalue.*`, which will be expanded to a list of the elements of the row value, just as occurs when the `.*` syntax is used at the top level of a `SELECT` list (see Section 8.16.5). For example, if table `t` has columns `f1` and `f2`, these are the same:

```sql
SELECT ROW(t.*, 42) FROM t;
SELECT ROW(t.f1, t.f2, 42) FROM t;
```

More expressions (separated by commas) for the row field values, and finally a right parenthesis. For example:

```sql
SELECT ROW(1,2.5,'this is a test');
```

The key word `ROW` is optional when there is more than one expression in the list.

A row constructor can include the syntax `rowvalue.*`, which will be expanded to a list of the elements of the row value, just as occurs when the `.*` syntax is used at the top level of a `SELECT` list (see Section 8.16.5). For example, if table `t` has columns `f1` and `f2`, these are the same:

```sql
SELECT ROW(t.*, 42) FROM t;
SELECT ROW(t.f1, t.f2, 42) FROM t;
```

**Note**

Before PostgreSQL 8.2, the `.*` syntax was not expanded in row constructors, so that writing `ROW(t.*, 42)` created a two-field row whose first field was another row value. The new behavior is usually more useful. If you need the old behavior of nested row values, write the inner row value without `.*`, for instance `ROW(t, 42).

By default, the value created by a `ROW` expression is of an anonymous record type. If necessary, it can be cast to a named composite type — either the row type of a table, or a composite type created with `CREATE TYPE AS`. An explicit cast might be needed to avoid ambiguity. For example:

```sql
CREATE TABLE mytable(f1 int, f2 float, f3 text);

CREATE FUNCTION getf1(mytable) RETURNS int AS 'SELECT $1.f1'
LANGUAGE SQL;

-- No cast needed since only one getf1() exists
SELECT getf1(ROW(1,2.5,'this is a test'));

getf1
-----
1
(1 row)

CREATE TYPE myrowtype AS (f1 int, f2 text, f3 numeric);

CREATE FUNCTION getf1(myrowtype) RETURNS int AS 'SELECT $1.f1'
LANGUAGE SQL;

-- Now we need a cast to indicate which function to call:
SELECT getf1(ROW(1,2.5,'this is a test'));
ERROR:  function getf1(record) is not unique

SELECT getf1(ROW(1,2.5,'this is a test')::mytable);
getf1
-----
1
(1 row)

SELECT getf1(CAST(ROW(1,2,5,'this is a test',2.5) AS myrowtype));
getf1
```

53
Row constructors can be used to build composite values to be stored in a composite-type table column, or to be passed to a function that accepts a composite parameter. Also, it is possible to compare two row values or test a row with IS NULL or IS NOT NULL, for example:

```
SELECT ROW(1,2.5,'this is a test') = ROW(1, 3, 'not the same');
SELECT ROW(table.*) IS NULL FROM table;  -- detect all-null rows
```

For more detail see Section 9.23. Row constructors can also be used in connection with subqueries, as discussed in Section 9.22.

### 4.2.14. Expression Evaluation Rules

The order of evaluation of subexpressions is not defined. In particular, the inputs of an operator or function are not necessarily evaluated left-to-right or in any other fixed order.

Furthermore, if the result of an expression can be determined by evaluating only some parts of it, then other subexpressions might not be evaluated at all. For instance, if one wrote:

```
SELECT true OR somefunc();
```

then `somefunc()` would (probably) not be called at all. The same would be the case if one wrote:

```
SELECT somefunc() OR true;
```

Note that this is not the same as the left-to-right “short-circuiting” of Boolean operators that is found in some programming languages.

As a consequence, it is unwise to use functions with side effects as part of complex expressions. It is particularly dangerous to rely on side effects or evaluation order in WHERE and HAVING clauses, since those clauses are extensively reprocessed as part of developing an execution plan. Boolean expressions (AND/OR/NOT combinations) in those clauses can be reorganized in any manner allowed by the laws of Boolean algebra.

When it is essential to force evaluation order, a CASE construct (see Section 9.17) can be used. For example, this is an untrustworthy way of trying to avoid division by zero in a WHERE clause:

```
SELECT ... WHERE x > 0 AND y/x > 1.5;
```

But this is safe:

```
SELECT ... WHERE CASE WHEN x > 0 THEN y/x > 1.5 ELSE false END;
```

A CASE construct used in this fashion will defeat optimization attempts, so it should only be done when necessary. (In this particular example, it would be better to sidestep the problem by writing `y > 1.5*x` instead.)

CASE is not a cure-all for such issues, however. One limitation of the technique illustrated above is that it does not prevent early evaluation of constant subexpressions. As described in Section 38.7, functions and operators marked IMMUTABLE can be evaluated when the query is planned rather than when it is executed. Thus for example
SELECT CASE WHEN x > 0 THEN x ELSE 1/0 END FROM tab;

is likely to result in a division-by-zero failure due to the planner trying to simplify the constant sub-expression, even if every row in the table has \( x > 0 \) so that the ELSE arm would never be entered at run time.

While that particular example might seem silly, related cases that don't obviously involve constants can occur in queries executed within functions, since the values of function arguments and local variables can be inserted into queries as constants for planning purposes. Within PL/pgSQL functions, for example, using an IF-THEN-ELSE statement to protect a risky computation is much safer than just nesting it in a CASE expression.

Another limitation of the same kind is that a CASE cannot prevent evaluation of an aggregate expression contained within it, because aggregate expressions are computed before other expressions in a SELECT list or HAVING clause are considered. For example, the following query can cause a division-by-zero error despite seemingly having protected against it:

```sql
SELECT CASE WHEN min(employees) > 0 THEN avg(expenses / employees) END FROM departments;
```

The `min()` and `avg()` aggregates are computed concurrently over all the input rows, so if any row has `employees` equal to zero, the division-by-zero error will occur before there is any opportunity to test the result of `min()`. Instead, use a WHERE or FILTER clause to prevent problematic input rows from reaching an aggregate function in the first place.

### 4.3. Calling Functions

PostgreSQL allows functions that have named parameters to be called using either positional or named notation. Named notation is especially useful for functions that have a large number of parameters, since it makes the associations between parameters and actual arguments more explicit and reliable. In positional notation, a function call is written with its argument values in the same order as they are defined in the function declaration. In named notation, the arguments are matched to the function parameters by name and can be written in any order. For each notation, also consider the effect of function argument types, documented in Section 10.3.

In either notation, parameters that have default values given in the function declaration need not be written in the call at all. But this is particularly useful in named notation, since any combination of parameters can be omitted; while in positional notation parameters can only be omitted from right to left.

PostgreSQL also supports mixed notation, which combines positional and named notation. In this case, positional parameters are written first and named parameters appear after them.

The following examples will illustrate the usage of all three notations, using the following function definition:

```sql
CREATE FUNCTION concat_lower_or_upper(a text, b text, uppercase boolean DEFAULT false) RETURNS text AS $$
SELECT CASE WHEN $3 THEN UPPER($1 || ' ' || $2)
ELSE LOWER($1 || ' ' || $2)
$$
```

```sql
SELECT CASE WHEN min(employees) > 0 THEN avg(expenses / employees) END FROM departments;
```
Function `concat_lower_or_upper` has two mandatory parameters, `a` and `b`. Additionally there is one optional parameter `uppercase` which defaults to `false`. The `a` and `b` inputs will be concatenated, and forced to either upper or lower case depending on the `uppercase` parameter. The remaining details of this function definition are not important here (see Chapter 38 for more information).

### 4.3.1. Using Positional Notation

Positional notation is the traditional mechanism for passing arguments to functions in PostgreSQL. An example is:

```sql
SELECT concat_lower_or_upper('Hello', 'World', true);
```

```
HELLO WORLD
(1 row)
```

All arguments are specified in order. The result is upper case since `uppercase` is specified as `true`. Another example is:

```sql
SELECT concat_lower_or_upper('Hello', 'World');
```

```
hello world
(1 row)
```

Here, the `uppercase` parameter is omitted, so it receives its default value of `false`, resulting in lower case output. In positional notation, arguments can be omitted from right to left so long as they have defaults.

### 4.3.2. Using Named Notation

In named notation, each argument's name is specified using `=>` to separate it from the argument expression. For example:

```sql
SELECT concat_lower_or_upper(a => 'Hello', b => 'World');
```

```
hello world
(1 row)
```

Again, the argument `uppercase` was omitted so it is set to `false` implicitly. One advantage of using named notation is that the arguments may be specified in any order, for example:

```sql
SELECT concat_lower_or_upper(a => 'Hello', b => 'World', uppercase => true);
```

```
HELLO WORLD
(1 row)
```
SELECT concat_lower_or_upper(a => 'Hello', uppercase => true, b => 'World');
concat_lower_or_upper
-----------------------
HELLO WORLD
(1 row)

An older syntax based on ":=" is supported for backward compatibility:

SELECT concat_lower_or_upper(a := 'Hello', uppercase := true, b := 'World');
concat_lower_or_upper
-----------------------
HELLO WORLD
(1 row)

4.3.3. Using Mixed Notation

The mixed notation combines positional and named notation. However, as already mentioned, named arguments cannot precede positional arguments. For example:

SELECT concat_lower_or_upper('Hello', 'World', uppercase => true);
concat_lower_or_upper
-----------------------
HELLO WORLD
(1 row)

In the above query, the arguments a and b are specified positionally, while uppercase is specified by name. In this example, that adds little except documentation. With a more complex function having numerous parameters that have default values, named or mixed notation can save a great deal of writing and reduce chances for error.

**Note**

Named and mixed call notations currently cannot be used when calling an aggregate function (but they do work when an aggregate function is used as a window function).
Chapter 5. Data Definition

This chapter covers how one creates the database structures that will hold one's data. In a relational database, the raw data is stored in tables, so the majority of this chapter is devoted to explaining how tables are created and modified and what features are available to control what data is stored in the tables. Subsequently, we discuss how tables can be organized into schemas, and how privileges can be assigned to tables. Finally, we will briefly look at other features that affect the data storage, such as inheritance, table partitioning, views, functions, and triggers.

5.1. Table Basics

A table in a relational database is much like a table on paper: It consists of rows and columns. The number and order of the columns is fixed, and each column has a name. The number of rows is variable — it reflects how much data is stored at a given moment. SQL does not make any guarantees about the order of the rows in a table. When a table is read, the rows will appear in an unspecified order, unless sorting is explicitly requested. This is covered in Chapter 7. Furthermore, SQL does not assign unique identifiers to rows, so it is possible to have several completely identical rows in a table. This is a consequence of the mathematical model that underlies SQL but is usually not desirable. Later in this chapter we will see how to deal with this issue.

Each column has a data type. The data type constrains the set of possible values that can be assigned to a column and assigns semantics to the data stored in the column so that it can be used for computations. For instance, a column declared to be of a numerical type will not accept arbitrary text strings, and the data stored in such a column can be used for mathematical computations. By contrast, a column declared to be of a character string type will accept almost any kind of data but it does not lend itself to mathematical calculations, although other operations such as string concatenation are available.

PostgreSQL includes a sizable set of built-in data types that fit many applications. Users can also define their own data types. Most built-in data types have obvious names and semantics, so we defer a detailed explanation to Chapter 8. Some of the frequently used data types are integer for whole numbers, numeric for possibly fractional numbers, text for character strings, date for dates, time for time-of-day values, and timestamp for values containing both date and time.

To create a table, you use the aptly named CREATE TABLE command. In this command you specify at least a name for the new table, the names of the columns and the data type of each column. For example:

```
CREATE TABLE my_first_table (
    first_column text,
    second_column integer
);
```

This creates a table named my_first_table with two columns. The first column is named first_column and has a data type of text; the second column has the name second_column and the type integer. The table and column names follow the identifier syntax explained in Section 4.1.1. The type names are usually also identifiers, but there are some exceptions. Note that the column list is comma-separated and surrounded by parentheses.

Of course, the previous example was heavily contrived. Normally, you would give names to your tables and columns that convey what kind of data they store. So let's look at a more realistic example:

```
CREATE TABLE products (
    product_no integer,
    name text,
```
Data Definition

    price numeric

);  

(The numeric type can store fractional components, as would be typical of monetary amounts.)

Tip

When you create many interrelated tables it is wise to choose a consistent naming pattern for the tables and columns. For instance, there is a choice of using singular or plural nouns for table names, both of which are favored by some theorist or other.

There is a limit on how many columns a table can contain. Depending on the column types, it is between 250 and 1600. However, defining a table with anywhere near this many columns is highly unusual and often a questionable design.

If you no longer need a table, you can remove it using the DROP TABLE command. For example:

    DROP TABLE my_first_table;
    DROP TABLE products;

Attempting to drop a table that does not exist is an error. Nevertheless, it is common in SQL script files to unconditionally try to drop each table before creating it, ignoring any error messages, so that the script works whether or not the table exists. (If you like, you can use the DROP TABLE IF EXISTS variant to avoid the error messages, but this is not standard SQL.)

If you need to modify a table that already exists, see Section 5.5 later in this chapter.

With the tools discussed so far you can create fully functional tables. The remainder of this chapter is concerned with adding features to the table definition to ensure data integrity, security, or convenience. If you are eager to fill your tables with data now you can skip ahead to Chapter 6 and read the rest of this chapter later.

5.2. Default Values

A column can be assigned a default value. When a new row is created and no values are specified for some of the columns, those columns will be filled with their respective default values. A data manipulation command can also request explicitly that a column be set to its default value, without having to know what that value is. (Details about data manipulation commands are in Chapter 6.)

If no default value is declared explicitly, the default value is the null value. This usually makes sense because a null value can be considered to represent unknown data.

In a table definition, default values are listed after the column data type. For example:

    CREATE TABLE products (  
        product_no integer,  
        name text,  
        price numeric DEFAULT 9.99  
    );

The default value can be an expression, which will be evaluated whenever the default value is inserted (not when the table is created). A common example is for a timestamp column to have a default of CURRENT_TIMESTAMP, so that it gets set to the time of row insertion. Another common example is generating a “serial number” for each row. In PostgreSQL this is typically done by something like:
CREATE TABLE products (  
    product_no integer DEFAULT nextval('products_product_no_seq'),  
    ...  
);  

where the \texttt{nextval()} function supplies successive values from a \textit{sequence object} (see Section 9.16). This arrangement is sufficiently common that there's a special shorthand for it:

CREATE TABLE products (  
    product_no SERIAL,  
    ...  
);  

The \texttt{SERIAL} shorthand is discussed further in Section 8.1.4.

5.3. Constraints

Data types are a way to limit the kind of data that can be stored in a table. For many applications, however, the constraint they provide is too coarse. For example, a column containing a product price should probably only accept positive values. But there is no standard data type that accepts only positive numbers. Another issue is that you might want to constrain column data with respect to other columns or rows. For example, in a table containing product information, there should be only one row for each product number.

To that end, SQL allows you to define constraints on columns and tables. Constraints give you as much control over the data in your tables as you wish. If a user attempts to store data in a column that would violate a constraint, an error is raised. This applies even if the value came from the default value definition.

5.3.1. Check Constraints

A check constraint is the most generic constraint type. It allows you to specify that the value in a certain column must satisfy a Boolean (truth-value) expression. For instance, to require positive product prices, you could use:

```sql  
CREATE TABLE products (  
    product_no integer,  
    name text,  
    price numeric \texttt{CHECK (price > 0)}  
);  
```

As you see, the constraint definition comes after the data type, just like default value definitions. Default values and constraints can be listed in any order. A check constraint consists of the key word \texttt{CHECK} followed by an expression in parentheses. The check constraint expression should involve the column thus constrained, otherwise the constraint would not make too much sense.

You can also give the constraint a separate name. This clarifies error messages and allows you to refer to the constraint when you need to change it. The syntax is:

```sql  
CREATE TABLE products (  
    product_no integer,  
    name text,  
    price numeric \texttt{CONSTRAINT positive_price CHECK (price > 0)}  
);  
```
So, to specify a named constraint, use the key word `CONSTRAINT` followed by an identifier followed by the constraint definition. (If you don't specify a constraint name in this way, the system chooses a name for you.)

A check constraint can also refer to several columns. Say you store a regular price and a discounted price, and you want to ensure that the discounted price is lower than the regular price:

```sql
CREATE TABLE products (
    product_no integer,
    name text,
    price numeric CHECK (price > 0),
    discounted_price numeric CHECK (discounted_price > 0),
    CHECK (price > discounted_price)
);
```

The first two constraints should look familiar. The third one uses a new syntax. It is not attached to a particular column, instead it appears as a separate item in the comma-separated column list. Column definitions and these constraint definitions can be listed in mixed order.

We say that the first two constraints are column constraints, whereas the third one is a table constraint because it is written separately from any one column definition. Column constraints can also be written as table constraints, while the reverse is not necessarily possible, since a column constraint is supposed to refer to only the column it is attached to. (PostgreSQL doesn't enforce that rule, but you should follow it if you want your table definitions to work with other database systems.) The above example could also be written as:

```sql
CREATE TABLE products (
    product_no integer,
    name text,
    price numeric,
    CHECK (price > 0),
    discounted_price numeric,
    CHECK (discounted_price > 0),
    CHECK (price > discounted_price)
);
```

or even:

```sql
CREATE TABLE products (
    product_no integer,
    name text,
    price numeric CHECK (price > 0),
    discounted_price numeric,
    CHECK (discounted_price > 0 AND price > discounted_price)
);
```

It's a matter of taste.

Names can be assigned to table constraints in the same way as column constraints:

```sql
CREATE TABLE products (
    product_no integer,
    name text,
    price numeric,
    CHECK (price > 0),
    discounted_price numeric,
    CHECK (discounted_price > 0 AND price > discounted_price)
);
```
CHECK (price > 0),
discounted_price numeric,
CHECK (discounted_price > 0),
CONSTRAINT valid_discount CHECK (price > discounted_price));

It should be noted that a check constraint is satisfied if the check expression evaluates to true or the null value. Since most expressions will evaluate to the null value if any operand is null, they will not prevent null values in the constrained columns. To ensure that a column does not contain null values, the not-null constraint described in the next section can be used.

5.3.2. Not-Null Constraints

A not-null constraint simply specifies that a column must not assume the null value. A syntax example:

CREATE TABLE products (  
    product_no integer NOT NULL,
    name text NOT NULL,
    price numeric
);

A not-null constraint is always written as a column constraint. A not-null constraint is functionally equivalent to creating a check constraint CHECK (column_name IS NOT NULL), but in PostgreSQL creating an explicit not-null constraint is more efficient. The drawback is that you cannot give explicit names to not-null constraints created this way.

Of course, a column can have more than one constraint. Just write the constraints one after another:

CREATE TABLE products (  
    product_no integer NOT NULL,
    name text NOT NULL,
    price numeric NOT NULL CHECK (price > 0)
);

The order doesn't matter. It does not necessarily determine in which order the constraints are checked.

The NOT NULL constraint has an inverse: the NULL constraint. This does not mean that the column must be null, which would surely be useless. Instead, this simply selects the default behavior that the column might be null. The NULL constraint is not present in the SQL standard and should not be used in portable applications. (It was only added to PostgreSQL to be compatible with some other database systems.) Some users, however, like it because it makes it easy to toggle the constraint in a script file. For example, you could start with:

CREATE TABLE products (  
    product_no integer NULL,
    name text NULL,
    price numeric NULL
);

and then insert the NOT key word where desired.

Tip
In most database designs the majority of columns should be marked not null.
5.3.3. Unique Constraints

Unique constraints ensure that the data contained in a column, or a group of columns, is unique among all the rows in the table. The syntax is:

```sql
CREATE TABLE products (
    product_no integer UNIQUE,
    name text,
    price numeric
);
```

when written as a column constraint, and:

```sql
CREATE TABLE products (
    product_no integer,
    name text,
    price numeric,
    UNIQUE (product_no)
);
```

when written as a table constraint.

To define a unique constraint for a group of columns, write it as a table constraint with the column names separated by commas:

```sql
CREATE TABLE example (
    a integer,
    b integer,
    c integer,
    UNIQUE (a, c)
);
```

This specifies that the combination of values in the indicated columns is unique across the whole table, though any one of the columns need not be (and ordinarily isn’t) unique.

You can assign your own name for a unique constraint, in the usual way:

```sql
CREATE TABLE products (
    product_no integer CONSTRAINT must_be_different UNIQUE,
    name text,
    price numeric
);
```

Adding a unique constraint will automatically create a unique B-tree index on the column or group of columns listed in the constraint. A uniqueness restriction covering only some rows cannot be written as a unique constraint, but it is possible to enforce such a restriction by creating a unique partial index.

In general, a unique constraint is violated if there is more than one row in the table where the values of all of the columns included in the constraint are equal. However, two null values are never considered equal in this comparison. That means even in the presence of a unique constraint it is possible to store duplicate rows that contain a null value in at least one of the constrained columns. This behavior conforms to the SQL standard, but we have heard that other SQL databases might not follow this rule. So be careful when developing applications that are intended to be portable.

5.3.4. Primary Keys
A primary key constraint indicates that a column, or group of columns, can be used as a unique identifier for rows in the table. This requires that the values be both unique and not null. So, the following two table definitions accept the same data:

```
CREATE TABLE products (
    product_no integer UNIQUE NOT NULL,
    name text,
    price numeric
);
```

```
CREATE TABLE products (
    product_no integer PRIMARY KEY,
    name text,
    price numeric
);
```

Primary keys can span more than one column; the syntax is similar to unique constraints:

```
CREATE TABLE example (
    a integer,
    b integer,
    c integer,
    PRIMARY KEY (a, c)
);
```

Adding a primary key will automatically create a unique B-tree index on the column or group of columns listed in the primary key, and will force the column(s) to be marked \texttt{NOT NULL}.

A table can have at most one primary key. (There can be any number of unique and not-null constraints, which are functionally almost the same thing, but only one can be identified as the primary key.) Relational database theory dictates that every table must have a primary key. This rule is not enforced by PostgreSQL, but it is usually best to follow it.

Primary keys are useful both for documentation purposes and for client applications. For example, a GUI application that allows modifying row values probably needs to know the primary key of a table to be able to identify rows uniquely. There are also various ways in which the database system makes use of a primary key if one has been declared; for example, the primary key defines the default target column(s) for foreign keys referencing its table.

### 5.3.5. Foreign Keys

A foreign key constraint specifies that the values in a column (or a group of columns) must match the values appearing in some row of another table. We say this maintains the \textit{referential integrity} between two related tables.

Say you have the product table that we have used several times already:

```
CREATE TABLE products (
    product_no integer PRIMARY KEY,
    name text,
    price numeric
);
```

Let's also assume you have a table storing orders of those products. We want to ensure that the orders table only contains orders of products that actually exist. So we define a foreign key constraint in the orders table that references the products table:
CREATE TABLE orders (  
  order_id integer PRIMARY KEY,  
  product_no integer REFERENCES products (product_no),  
  quantity integer  
);  

Now it is impossible to create orders with non-NULL product_no entries that do not appear in the products table.

We say that in this situation the orders table is the referencing table and the products table is the referenced table. Similarly, there are referencing and referenced columns.

You can also shorten the above command to:

CREATE TABLE orders (  
  order_id integer PRIMARY KEY,  
  product_no integer REFERENCES products,  
  quantity integer  
);  

because in absence of a column list the primary key of the referenced table is used as the referenced column(s).

A foreign key can also constrain and reference a group of columns. As usual, it then needs to be written in table constraint form. Here is a contrived syntax example:

CREATE TABLE t1 (  
  a integer PRIMARY KEY,  
  b integer,  
  c integer,  
  FOREIGN KEY (b, c) REFERENCES other_table (c1, c2)  
);  

Of course, the number and type of the constrained columns need to match the number and type of the referenced columns.

You can assign your own name for a foreign key constraint, in the usual way.

A table can have more than one foreign key constraint. This is used to implement many-to-many relationships between tables. Say you have tables about products and orders, but now you want to allow one order to contain possibly many products (which the structure above did not allow). You could use this table structure:

CREATE TABLE products (  
  product_no integer PRIMARY KEY,  
  name text,  
  price numeric  
);  

CREATE TABLE orders (  
  order_id integer PRIMARY KEY,  
  shipping_address text,  
  ...  
);  

CREATE TABLE order_items (  
  order_id integer,  
  product_no integer REFERENCES products (product_no),  
  quantity integer  
);
Data Definition

```
product_no integer REFERENCES products,
order_id integer REFERENCES orders,
quantity integer,
PRIMARY KEY (product_no, order_id)
```

Notice that the primary key overlaps with the foreign keys in the last table.

We know that the foreign keys disallow creation of orders that do not relate to any products. But what if a product is removed after an order is created that references it? SQL allows you to handle that as well. Intuitively, we have a few options:

- Disallow deleting a referenced product
- Delete the orders as well
- Something else?

To illustrate this, let's implement the following policy on the many-to-many relationship example above: when someone wants to remove a product that is still referenced by an order (via order_items), we disallow it. If someone removes an order, the order items are removed as well:

```
CREATE TABLE products (  
    product_no integer PRIMARY KEY,  
    name text,  
    price numeric  
);

CREATE TABLE orders (  
    order_id integer PRIMARY KEY,  
    shipping_address text,  
    ...  
);

CREATE TABLE order_items (  
    product_no integer REFERENCES products  
        ON DELETE RESTRICT,  
    order_id integer REFERENCES orders  
        ON DELETE CASCADE,  
    quantity integer,  
    PRIMARY KEY (product_no, order_id)  
);
```

Restricting and cascading deletes are the two most common options. RESTRICT prevents deletion of a referenced row. NO ACTION means that if any referencing rows still exist when the constraint is checked, an error is raised; this is the default behavior if you do not specify anything. (The essential difference between these two choices is that NO ACTION allows the check to be deferred until later in the transaction, whereas RESTRICT does not.) CASCADE specifies that when a referenced row is deleted, row(s) referencing it should be automatically deleted as well. There are two other options: SET NULL and SET DEFAULT. These cause the referencing column(s) in the referencing row(s) to be set to nulls or their default values, respectively, when the referenced row is deleted. Note that these do not excuse you from observing any constraints. For example, if an action specifies SET DEFAULT but the default value would not satisfy the foreign key constraint, the operation will fail.

Analogous to ON DELETE there is also ON UPDATE which is invoked when a referenced column is changed (updated). The possible actions are the same. In this case, CASCADE means that the updated values of the referenced column(s) should be copied into the referencing row(s).

Normally, a referencing row need not satisfy the foreign key constraint if any of its referencing columns are null. If MATCH FULL is added to the foreign key declaration, a referencing row escapes satisfying the constraint only if all its referencing columns are null (so a mix of null and non-null
values is guaranteed to fail a MATCH FULL constraint). If you don't want referencing rows to be able to avoid satisfying the foreign key constraint, declare the referencing column(s) as NOT NULL.

A foreign key must reference columns that either are a primary key or form a unique constraint. This means that the referenced columns always have an index (the one underlying the primary key or unique constraint); so checks on whether a referencing row has a match will be efficient. Since a DELETE of a row from the referenced table or an UPDATE of a referenced column will require a scan of the referencing table for rows matching the old value, it is often a good idea to index the referencing columns too. Because this is not always needed, and there are many choices available on how to index, declaration of a foreign key constraint does not automatically create an index on the referencing columns.

More information about updating and deleting data is in Chapter 6. Also see the description of foreign key constraint syntax in the reference documentation for CREATE TABLE.

5.3.6. Exclusion Constraints

Exclusion constraints ensure that if any two rows are compared on the specified columns or expressions using the specified operators, at least one of these operator comparisons will return false or null. The syntax is:

```sql
CREATE TABLE circles (  
c circle,  
    EXCLUDE USING gist (c WITH &&)  
);  
```

See also CREATE TABLE ... CONSTRAINT ... EXCLUDE for details.

Adding an exclusion constraint will automatically create an index of the type specified in the constraint declaration.

5.4. System Columns

Every table has several system columns that are implicitly defined by the system. Therefore, these names cannot be used as names of user-defined columns. (Note that these restrictions are separate from whether the name is a key word or not; quoting a name will not allow you to escape these restrictions.) You do not really need to be concerned about these columns; just know they exist.

**oid**

The object identifier (object ID) of a row. This column is only present if the table was created using WITH OIDS, or if the default_with_oids configuration variable was set at the time. This column is of type oid (same name as the column); see Section 8.19 for more information about the type.

**tableoid**

The OID of the table containing this row. This column is particularly handy for queries that select from inheritance hierarchies (see Section 5.9), since without it, it's difficult to tell which individual table a row came from. The tableoid can be joined against the oid column of pg_class to obtain the table name.

**xmin**

The identity (transaction ID) of the inserting transaction for this row version. (A row version is an individual state of a row; each update of a row creates a new row version for the same logical row.)

**cmin**

67
The command identifier (starting at zero) within the inserting transaction.

\( xmax \)

The identity (transaction ID) of the deleting transaction, or zero for an undeleted row version. It is possible for this column to be nonzero in a visible row version. That usually indicates that the deleting transaction hasn't committed yet, or that an attempted deletion was rolled back.

\( cmax \)

The command identifier within the deleting transaction, or zero.

\( ctid \)

The physical location of the row version within its table. Note that although the \( ctid \) can be used to locate the row version very quickly, a row's \( ctid \) will change if it is updated or moved by VACUUM FULL. Therefore \( ctid \) is useless as a long-term row identifier. The OID, or even better a user-defined serial number, should be used to identify logical rows.

OIDs are 32-bit quantities and are assigned from a single cluster-wide counter. In a large or long-lived database, it is possible for the counter to wrap around. Hence, it is bad practice to assume that OIDS are unique, unless you take steps to ensure that this is the case. If you need to identify the rows in a table, using a sequence generator is strongly recommended. However, OIDS can be used as well, provided that a few additional precautions are taken:

- A unique constraint should be created on the OID column of each table for which the OID will be used to identify rows. When such a unique constraint (or unique index) exists, the system takes care not to generate an OID matching an already-existing row. (Of course, this is only possible if the table contains fewer than \( 2^{32} \) (4 billion) rows, and in practice the table size had better be much less than that, or performance might suffer.)

- OIDS should never be assumed to be unique across tables; use the combination of \( \text{tableoid} \) and row OID if you need a database-wide identifier.

- Of course, the tables in question must be created \( \text{WITH OIDS} \). As of PostgreSQL 8.1, \( \text{WITHOUT OIDS} \) is the default.

Transaction identifiers are also 32-bit quantities. In a long-lived database it is possible for transaction IDs to wrap around. This is not a fatal problem given appropriate maintenance procedures; see Chapter 24 for details. It is unwise, however, to depend on the uniqueness of transaction IDs over the long term (more than one billion transactions).

Command identifiers are also 32-bit quantities. This creates a hard limit of \( 2^{32} \) (4 billion) SQL commands within a single transaction. In practice this limit is not a problem — note that the limit is on the number of SQL commands, not the number of rows processed. Also, only commands that actually modify the database contents will consume a command identifier.

## 5.5. Modifying Tables

When you create a table and you realize that you made a mistake, or the requirements of the application change, you can drop the table and create it again. But this is not a convenient option if the table is already filled with data, or if the table is referenced by other database objects (for instance a foreign key constraint). Therefore PostgreSQL provides a family of commands to make modifications to existing tables. Note that this is conceptually distinct from altering the data contained in the table: here we are interested in altering the definition, or structure, of the table.

You can:

- Add columns
- Remove columns
• Add constraints
• Remove constraints
• Change default values
• Change column data types
• Rename columns
• Rename tables

All these actions are performed using the ALTER TABLE command, whose reference page contains details beyond those given here.

5.5.1. Adding a Column

To add a column, use a command like:

```
ALTER TABLE products ADD COLUMN description text;
```

The new column is initially filled with whatever default value is given (null if you don't specify a DEFAULT clause).

You can also define constraints on the column at the same time, using the usual syntax:

```
ALTER TABLE products ADD COLUMN description text CHECK (description <> '');
```

In fact all the options that can be applied to a column description in CREATE TABLE can be used here. Keep in mind however that the default value must satisfy the given constraints, or the ADD will fail. Alternatively, you can add constraints later (see below) after you've filled in the new column correctly.

**Tip**

Adding a column with a default requires updating each row of the table (to store the new column value). However, if no default is specified, PostgreSQL is able to avoid the physical update. So if you intend to fill the column with mostly nondefault values, it's best to add the column with no default, insert the correct values using UPDATE, and then add any desired default as described below.

5.5.2. Removing a Column

To remove a column, use a command like:

```
ALTER TABLE products DROP COLUMN description;
```

Whatever data was in the column disappears. Table constraints involving the column are dropped, too. However, if the column is referenced by a foreign key constraint of another table, PostgreSQL will not silently drop that constraint. You can authorize dropping everything that depends on the column by adding CASCADE:

```
ALTER TABLE products DROP COLUMN description CASCADE;
```

See Section 5.13 for a description of the general mechanism behind this.

5.5.3. Adding a Constraint
To add a constraint, the table constraint syntax is used. For example:

```sql
ALTER TABLE products ADD CHECK (name <> '');
ALTER TABLE products ADD CONSTRAINT some_name UNIQUE (product_no);
ALTER TABLE products ADD FOREIGN KEY (product_group_id) REFERENCES product_groups;
```

To add a not-null constraint, which cannot be written as a table constraint, use this syntax:

```sql
ALTER TABLE products ALTER COLUMN product_no SET NOT NULL;
```

The constraint will be checked immediately, so the table data must satisfy the constraint before it can be added.

### 5.5.4. Removing a Constraint

To remove a constraint you need to know its name. If you gave it a name then that’s easy. Otherwise the system assigned a generated name, which you need to find out. The `psql` command `\d table-name` can be helpful here; other interfaces might also provide a way to inspect table details. Then the command is:

```sql
ALTER TABLE products DROP CONSTRAINT some_name;
```

(If you are dealing with a generated constraint name like `$2`, don’t forget that you’ll need to double-quote it to make it a valid identifier.)

As with dropping a column, you need to add `CASCADE` if you want to drop a constraint that something else depends on. An example is that a foreign key constraint depends on a unique or primary key constraint on the referenced column(s).

This works the same for all constraint types except not-null constraints. To drop a not null constraint use:

```sql
ALTER TABLE products ALTER COLUMN product_no DROP NOT NULL;
```

(Recall that not-null constraints do not have names.)

### 5.5.5. Changing a Column’s Default Value

To set a new default for a column, use a command like:

```sql
ALTER TABLE products ALTER COLUMN price SET DEFAULT 7.77;
```

Note that this doesn’t affect any existing rows in the table, it just changes the default for future `INSERT` commands.

To remove any default value, use:

```sql
ALTER TABLE products ALTER COLUMN price DROP DEFAULT;
```

This is effectively the same as setting the default to null. As a consequence, it is not an error to drop a default where one hadn’t been defined, because the default is implicitly the null value.

### 5.5.6. Changing a Column's Data Type
To convert a column to a different data type, use a command like:

```
ALTER TABLE products ALTER COLUMN price TYPE numeric(10,2);
```

This will succeed only if each existing entry in the column can be converted to the new type by an implicit cast. If a more complex conversion is needed, you can add a `USING` clause that specifies how to compute the new values from the old.

PostgreSQL will attempt to convert the column's default value (if any) to the new type, as well as any constraints that involve the column. But these conversions might fail, or might produce surprising results. It's often best to drop any constraints on the column before altering its type, and then add back suitably modified constraints afterwards.

### 5.5.7. Renaming a Column

To rename a column:

```
ALTER TABLE products RENAME COLUMN product_no TO product_number;
```

### 5.5.8. Renaming a Table

To rename a table:

```
ALTER TABLE products RENAME TO items;
```

### 5.6. Privileges

When an object is created, it is assigned an owner. The owner is normally the role that executed the creation statement. For most kinds of objects, the initial state is that only the owner (or a superuser) can do anything with the object. To allow other roles to use it, privileges must be granted.

There are different kinds of privileges: `SELECT`, `INSERT`, `UPDATE`, `DELETE`, `TRUNCATE`, `REFERENCES`, `TRIGGER`, `CREATE`, `CONNECT`, `TEMPORARY`, `EXECUTE`, and `USAGE`. The privileges applicable to a particular object vary depending on the object's type (table, function, etc). For complete information on the different types of privileges supported by PostgreSQL, refer to the GRANT reference page. The following sections and chapters will also show you how those privileges are used.

The right to modify or destroy an object is always the privilege of the owner only.

An object can be assigned to a new owner with an `ALTER` command of the appropriate kind for the object, e.g. `ALTER TABLE`. Superusers can always do this; ordinary roles can only do it if they are both the current owner of the object (or a member of the owning role) and a member of the new owning role.

To assign privileges, the `GRANT` command is used. For example, if `joe` is an existing role, and `accounts` is an existing table, the privilege to update the table can be granted with:

```
GRANT UPDATE ON accounts TO joe;
```

Writing `ALL` in place of a specific privilege grants all privileges that are relevant for the object type.

The special “role” name `PUBLIC` can be used to grant a privilege to every role on the system. Also, “group” roles can be set up to help manage privileges when there are many users of a database — for details see Chapter 21.
To revoke a privilege, use the fittingly named `REVOKE` command:

```sql
REVOKE ALL ON accounts FROM PUBLIC;
```

The special privileges of the object owner (i.e., the right to do DROP, GRANT, REVOKE, etc.) are always implicit in being the owner, and cannot be granted or revoked. But the object owner can choose to revoke their own ordinary privileges, for example to make a table read-only for themselves as well as others.

Ordinarily, only the object's owner (or a superuser) can grant or revoke privileges on an object. However, it is possible to grant a privilege “with grant option”, which gives the recipient the right to grant it in turn to others. If the grant option is subsequently revoked then all who received the privilege from that recipient (directly or through a chain of grants) will lose the privilege. For details see the GRANT and REVOKE reference pages.

## 5.7. Row Security Policies

In addition to the SQL-standard privilege system available through GRANT, tables can have row security policies that restrict, on a per-user basis, which rows can be returned by normal queries or inserted, updated, or deleted by data modification commands. This feature is also known as Row-Level Security. By default, tables do not have any policies, so that if a user has access privileges to a table according to the SQL privilege system, all rows within it are equally available for querying or updating.

When row security is enabled on a table (with `ALTER TABLE ... ENABLE ROW LEVEL SECURITY`), all normal access to the table for selecting rows or modifying rows must be allowed by a row security policy. (However, the table's owner is typically not subject to row security policies.) If no policy exists for the table, a default-deny policy is used, meaning that no rows are visible or can be modified. Operations that apply to the whole table, such as `TRUNCATE` and `REFERENCES`, are not subject to row security.

Row security policies can be specific to commands, or to roles, or to both. A policy can be specified to apply to all commands, or to `SELECT`, `INSERT`, `UPDATE`, or `DELETE`. Multiple roles can be assigned to a given policy, and normal role membership and inheritance rules apply.

To specify which rows are visible or modifiable according to a policy, an expression is required that returns a Boolean result. This expression will be evaluated for each row prior to any conditions or functions coming from the user's query. (The only exceptions to this rule are leakproof functions, which are guaranteed not to leak information; the optimizer may choose to apply such functions ahead of the row-security check.) Rows for which the expression does not return `true` will not be processed.

Separate expressions may be specified to provide independent control over the rows which are visible and the rows which are allowed to be modified. Policy expressions are run as part of the query and with the privileges of the user running the query, although security-definer functions can be used to access data not available to the calling user.

Superusers and roles with the `BYPASSRLS` attribute always bypass the row security system when accessing a table. Table owners normally bypass row security as well, though a table owner can choose to be subject to row security with `ALTER TABLE ... FORCE ROW LEVEL SECURITY`.

Enabling and disabling row security, as well as adding policies to a table, is always the privilege of the table owner only.

Policies are created using the `CREATE POLICY` command, altered using the `ALTER POLICY` command, and dropped using the `DROP POLICY` command. To enable and disable row security for a given table, use the `ALTER TABLE` command.

Each policy has a name and multiple policies can be defined for a table. As policies are table-specific, each policy for a table must have a unique name. Different tables may have policies with the same name.
When multiple policies apply to a given query, they are combined using either OR (for permissive policies, which are the default) or using AND (for restrictive policies). This is similar to the rule that a given role has the privileges of all roles that they are a member of. Permissive vs. restrictive policies are discussed further below.

As a simple example, here is how to create a policy on the `accounts` relation to allow only members of the `managers` role to access rows, and only rows of their accounts:

```sql
CREATE TABLE accounts (manager text, company text, contact_email text);
ALTER TABLE accounts ENABLE ROW LEVEL SECURITY;
CREATE POLICY account_managers ON accounts TO managers
    USING (manager = current_user);
```

The policy above implicitly provides a `WITH CHECK` clause identical to its `USING` clause, so that the constraint applies both to rows selected by a command (so a manager cannot `SELECT`, `UPDATE`, or `DELETE` existing rows belonging to a different manager) and to rows modified by a command (so rows belonging to a different manager cannot be created via `INSERT` or `UPDATE`).

If no role is specified, or the special user name `PUBLIC` is used, then the policy applies to all users on the system. To allow all users to access only their own row in a `users` table, a simple policy can be used:

```sql
CREATE POLICY user_policy ON users
    USING (user_name = current_user);
```

This works similarly to the previous example.

To use a different policy for rows that are being added to the table compared to those rows that are visible, multiple policies can be combined. This pair of policies would allow all users to view all rows in the `users` table, but only modify their own:

```sql
CREATE POLICY user_sel_policy ON users
    FOR SELECT
    USING (true);
CREATE POLICY user_mod_policy ON users
    USING (user_name = current_user);
```

In a `SELECT` command, these two policies are combined using OR, with the net effect being that all rows can be selected. In other command types, only the second policy applies, so that the effects are the same as before.

Row security can also be disabled with the `ALTER TABLE` command. Disabling row security does not remove any policies that are defined on the table; they are simply ignored. Then all rows in the table are visible and modifiable, subject to the standard SQL privileges system.

Below is a larger example of how this feature can be used in production environments. The table `passwd` emulates a Unix password file:

```sql
-- Simple passwd-file based example
CREATE TABLE passwd ( 
  user_name             text UNIQUE NOT NULL,
  pwhash                text,
  uid                   int  PRIMARY KEY,
```
CREATE ROLE admin; -- Administrator
CREATE ROLE bob; -- Normal user
CREATE ROLE alice; -- Normal user

-- Populate the table
INSERT INTO passwd VALUES
  ('admin','xxx',0,0,'Admin','111-222-3333',null,'/root','/bin/dash');
INSERT INTO passwd VALUES
  ('bob','xxx',1,1,'Bob','123-456-7890',null,'/home/bob','/bin/zsh');
INSERT INTO passwd VALUES
  ('alice','xxx',2,1,'Alice','098-765-4321',null,'/home/alice','/bin/zsh');

-- Be sure to enable row level security on the table
ALTER TABLE passwd ENABLE ROW LEVEL SECURITY;

-- Create policies
-- Administrator can see all rows and add any rows
CREATE POLICY admin_all ON passwd TO admin USING (true) WITH CHECK (true);
-- Normal users can view all rows
CREATE POLICY all_view ON passwd FOR SELECT USING (true);
-- Normal users can update their own records, but
-- limit which shells a normal user is allowed to set
CREATE POLICY user_mod ON passwd FOR UPDATE USING (current_user = user_name)
  WITH CHECK ( current_user = user_name AND shell IN ('/bin/bash','/bin/sh','/bin/dash','/bin/zsh','/bin/tcsh') );

-- Allow admin all normal rights
GRANT SELECT, INSERT, UPDATE, DELETE ON passwd TO admin;
-- Users only get select access on public columns
GRANT SELECT
  (user_name, uid, gid, real_name, home_phone, extra_info, home_dir, shell)
ON passwd TO public;
-- Allow users to update certain columns
GRANT UPDATE
  (pwhash, real_name, home_phone, extra_info, shell)
ON passwd TO public;

As with any security settings, it's important to test and ensure that the system is behaving as expected. Using the example above, this demonstrates that the permission system is working properly.
-- admin can view all rows and fields
postgres=> set role admin;
SET
postgres=> table passwd;
user_name | pwhash | uid | gid | real_name | home_phone |
extra_info | home_dir    |   shell
-----------+--------+-----+-----+-----------+--------------
| /root       |
admin     | xxx    |   0 |   0 | Admin     | 111-222-3333 |
| /bin/dash |
bob       | xxx    |   1 |   1 | Bob       | 123-456-7890 |
| /home/bob   |
alice     | xxx    |   2 |   1 | Alice     | 098-765-4321 |
| /home/alice | /bin/zsh |
(3 rows)

-- Test what Alice is able to do
postgres=> set role alice;
SET
postgres=> table passwd;
ERROR:  permission denied for relation passwd
postgres=> select
user_name,real_name,home_phone,extra_info,home_dir,shell from
passwd;
user_name | real_name |  home_phone  | extra_info | home_dir    |
shell
-----------+-----------+--------------+------------+-------------
| /root     |
admin     | Admin     | 111-222-3333 |            | /root       |
| /bin/dash |
bob       | Bob       | 123-456-7890 |            | /home/bob   |
| /bin/zsh |
alice     | Alice     | 098-765-4321 |            | /home/alice |
| /bin/zsh |
(3 rows)

postgres=> update passwd set user_name = 'joe';
ERROR:  permission denied for relation passwd
-- Alice is allowed to change her own real_name, but no others
postgres=> update passwd set real_name = 'Alice Doe';
UPDATE 1
postgres=> update passwd set real_name = 'John Doe' where user_name = 'admin';
UPDATE 0
postgres=> update passwd set shell = '/bin/xx';
ERROR:  new row violates WITH CHECK OPTION for "passwd"
postgres=> delete from passwd;
ERROR:  permission denied for relation passwd
postgres=> insert into passwd (user_name) values ('xxx');
ERROR:  permission denied for relation passwd
-- Alice can change her own password; RLS silently prevents updating other rows
postgres=> update passwd set pwhash = 'abc';
UPDATE 1

All of the policies constructed thus far have been permissive policies, meaning that when multiple policies are applied they are combined using the “OR” Boolean operator. While permissive policies can be constructed to only allow access to rows in the intended cases, it can be simpler to combine
permissive policies with restrictive policies (which the records must pass and which are combined using the “AND” Boolean operator). Building on the example above, we add a restrictive policy to require the administrator to be connected over a local Unix socket to access the records of the passwd table:

```
CREATE POLICY admin_local_only ON passwd AS RESTRICTIVE TO admin
    USING (pg_catalog.inet_client_addr() IS NULL);
```

We can then see that an administrator connecting over a network will not see any records, due to the restrictive policy:

```
=> SELECT current_user;
   current_user
       --------
        admin
      (1 row)

=> SELECT inet_client_addr();
   inet_client_addr
       ------------
          127.0.0.1
      (1 row)

=> SELECT current_user;
   current_user
       --------
        admin
      (1 row)

=> TABLE passwd;
user_name | pwhash | uid | gid | real_name | home_phone | extra_info | home_dir | shell
---------+--------+-----+-----+-----------+------------+-----------+---------+-------
         +-----------------------------+-----------
      (0 rows)

=> UPDATE passwd set pwhash = NULL;
   UPDATE 0
```

Referential integrity checks, such as unique or primary key constraints and foreign key references, always bypass row security to ensure that data integrity is maintained. Care must be taken when developing schemas and row level policies to avoid “covert channel” leaks of information through such referential integrity checks.

In some contexts it is important to be sure that row security is not being applied. For example, when taking a backup, it could be disastrous if row security silently caused some rows to be omitted from the backup. In such a situation, you can set the row_security configuration parameter to off. This does not in itself bypass row security; what it does is throw an error if any query’s results would get filtered by a policy. The reason for the error can then be investigated and fixed.

In the examples above, the policy expressions consider only the current values in the row to be accessed or updated. This is the simplest and best-performing case; when possible, it’s best to design row security applications to work this way. If it is necessary to consult other rows or other tables to make a policy decision, that can be accomplished using sub-SELECTs, or functions that contain SELECTs, in the policy expressions. Be aware however that such accesses can create race conditions that could allow information leakage if care is not taken. As an example, consider the following table design:
-- definition of privilege groups
CREATE TABLE groups (group_id int PRIMARY KEY,
    group_name text NOT NULL);

INSERT INTO groups VALUES
    (1, 'low'),
    (2, 'medium'),
    (5, 'high');

GRANT ALL ON groups TO alice; -- alice is the administrator
GRANT SELECT ON groups TO public;

-- definition of users' privilege levels
CREATE TABLE users (user_name text PRIMARY KEY,
    group_id int NOT NULL REFERENCES groups);

INSERT INTO users VALUES
    ('alice', 5),
    ('bob', 2),
    ('mallory', 2);

GRANT ALL ON users TO alice;
GRANT SELECT ON users TO public;

-- table holding the information to be protected
CREATE TABLE information (info text,
    group_id int NOT NULL REFERENCES groups);

INSERT INTO information VALUES
    ('barely secret', 1),
    ('slightly secret', 2),
    ('very secret', 5);

ALTER TABLE information ENABLE ROW LEVEL SECURITY;

-- a row should be visible to/updatable by users whose security
-- group_id is
-- greater than or equal to the row's group_id
CREATE POLICY fp_s ON information FOR SELECT
    USING (group_id <= (SELECT group_id FROM users WHERE user_name =
        current_user));
CREATE POLICY fp_u ON information FOR UPDATE
    USING (group_id <= (SELECT group_id FROM users WHERE user_name =
        current_user));

-- we rely only on RLS to protect the information table
GRANT ALL ON information TO public;

Now suppose that alice wishes to change the "slightly secret" information, but decides that mallory should not be trusted with the new content of that row, so she does:

BEGIN;
UPDATE users SET group_id = 1 WHERE user_name = 'mallory';
UPDATE information SET info = 'secret from mallory' WHERE group_id = 2;
COMMIT;
That looks safe; there is no window wherein mallory should be able to see the “secret from mallory” string. However, there is a race condition here. If mallory is concurrently doing, say,

```
SELECT * FROM information WHERE group_id = 2 FOR UPDATE;
```

and her transaction is in READ COMMITTED mode, it is possible for her to see “secret from mallory”. That happens if her transaction reaches the information row just after alice’s does. It blocks waiting for alice’s transaction to commit, then fetches the updated row contents thanks to the FOR UPDATE clause. However, it does not fetch an updated row for the implicit SELECT from users, because that sub-SELECT did not have FOR UPDATE; instead the users row is read with the snapshot taken at the start of the query. Therefore, the policy expression tests the old value of mallory’s privilege level and allows her to see the updated row.

There are several ways around this problem. One simple answer is to use SELECT ... FOR SHARE in sub-SELECTs in row security policies. However, that requires granting UPDATE privilege on the referenced table (here users) to the affected users, which might be undesirable. (But another row security policy could be applied to prevent them from actually exercising that privilege; or the sub-SELECT could be embedded into a security definer function.) Also, heavy concurrent use of row share locks on the referenced table could pose a performance problem, especially if updates of it are frequent. Another solution, practical if updates of the referenced table are infrequent, is to take an exclusive lock on the referenced table when updating it, so that no concurrent transactions could be examining old row values. Or one could just wait for all concurrent transactions to end after committing an update of the referenced table and before making changes that rely on the new security situation.

For additional details see CREATE POLICY and ALTER TABLE.

### 5.8. Schemas

A PostgreSQL database cluster contains one or more named databases. Users and groups of users are shared across the entire cluster, but no other data is shared across databases. Any given client connection to the server can access only the data in a single database, the one specified in the connection request.

Users of a cluster do not necessarily have the privilege to access every database in the cluster. Sharing of user names means that there cannot be different users named, say, joe in two databases in the same cluster; but the system can be configured to allow joe access to only some of the databases.

A database contains one or more named schemas, which in turn contain tables. Schemas also contain other kinds of named objects, including data types, functions, and operators. The same object name can be used in different schemas without conflict; for example, both schemal and myschema can contain tables named mytable. Unlike databases, schemas are not rigidly separated: a user can access objects in any of the schemas in the database they are connected to, if they have privileges to do so.

There are several reasons why one might want to use schemas:

- To allow many users to use one database without interfering with each other.
- To organize database objects into logical groups to make them more manageable.
- Third-party applications can be put into separate schemas so they do not collide with the names of other objects.

Schemas are analogous to directories at the operating system level, except that schemas cannot be nested.
5.8.1. Creating a Schema

To create a schema, use the CREATE SCHEMA command. Give the schema a name of your choice. For example:

```sql
CREATE SCHEMA myschema;
```

To create or access objects in a schema, write a qualified name consisting of the schema name and table name separated by a dot:

```sql
schema.table
```

This works anywhere a table name is expected, including the table modification commands and the data access commands discussed in the following chapters. (For brevity we will speak of tables only, but the same ideas apply to other kinds of named objects, such as types and functions.)

Actually, the even more general syntax

```sql
database.schema.table
```

can be used too, but at present this is just for pro forma compliance with the SQL standard. If you write a database name, it must be the same as the database you are connected to.

So to create a table in the new schema, use:

```sql
CREATE TABLE myschema.mytable {
  ...
};
```

To drop a schema if it’s empty (all objects in it have been dropped), use:

```sql
DROP SCHEMA myschema;
```

To drop a schema including all contained objects, use:

```sql
DROP SCHEMA myschema CASCADE;
```

See Section 5.13 for a description of the general mechanism behind this.

Often you will want to create a schema owned by someone else (since this is one of the ways to restrict the activities of your users to well-defined namespaces). The syntax for that is:

```sql
CREATE SCHEMA schema_name AUTHORIZATION user_name;
```

You can even omit the schema name, in which case the schema name will be the same as the user name. See Section 5.8.6 for how this can be useful.

Schema names beginning with `pg_` are reserved for system purposes and cannot be created by users.

5.8.2. The Public Schema
In the previous sections we created tables without specifying any schema names. By default such tables (and other objects) are automatically put into a schema named “public”. Every new database contains such a schema. Thus, the following are equivalent:

```
CREATE TABLE products ( ... );
```

and:

```
CREATE TABLE public.products ( ... );
```

### 5.8.3. The Schema Search Path

Qualified names are tedious to write, and it's often best not to wire a particular schema name into applications anyway. Therefore tables are often referred to by unqualified names, which consist of just the table name. The system determines which table is meant by following a search path, which is a list of schemas to look in. The first matching table in the search path is taken to be the one wanted. If there is no match in the search path, an error is reported, even if matching table names exist in other schemas in the database.

The ability to create like-named objects in different schemas complicates writing a query that references precisely the same objects every time. It also opens up the potential for users to change the behavior of other users' queries, maliciously or accidentally. Due to the prevalence of unqualified names in queries and their use in PostgreSQL internals, adding a schema to `search_path` effectively trusts all users having `CREATE` privilege on that schema. When you run an ordinary query, a malicious user able to create objects in a schema of your search path can take control and execute arbitrary SQL functions as though you executed them.

The first schema named in the search path is called the current schema. Aside from being the first schema searched, it is also the schema in which new tables will be created if the `CREATE TABLE` command does not specify a schema name.

To show the current search path, use the following command:

```
SHOW search_path;
```

In the default setup this returns:

```
search_path
-----------
"$user", public
```

The first element specifies that a schema with the same name as the current user is to be searched. If no such schema exists, the entry is ignored. The second element refers to the public schema that we have seen already.

The first schema in the search path that exists is the default location for creating new objects. That is the reason that by default objects are created in the public schema. When objects are referenced in any other context without schema qualification (table modification, data modification, or query commands) the search path is traversed until a matching object is found. Therefore, in the default configuration, any unqualified access again can only refer to the public schema.

To put our new schema in the path, we use:
SET search_path TO myschema, public;

(We omit the $user here because we have no immediate need for it.) And then we can access the table without schema qualification:

DROP TABLE mytable;

Also, since myschema is the first element in the path, new objects would by default be created in it.

We could also have written:

SET search_path TO myschema;

Then we no longer have access to the public schema without explicit qualification. There is nothing special about the public schema except that it exists by default. It can be dropped, too.

See also Section 9.25 for other ways to manipulate the schema search path.

The search path works in the same way for data type names, function names, and operator names as it does for table names. Data type and function names can be qualified in exactly the same way as table names. If you need to write a qualified operator name in an expression, there is a special provision: you must write

\[ \text{OPERATOR}\{(\text{schema}.\text{operator})\} \]

This is needed to avoid syntactic ambiguity. An example is:

\[ \text{SELECT 3 OPERATOR(pg_catalog.+)} 4; \]

In practice one usually relies on the search path for operators, so as not to have to write anything so ugly as that.

### 5.8.4. Schemas and Privileges

By default, users cannot access any objects in schemas they do not own. To allow that, the owner of the schema must grant the USAGE privilege on the schema. To allow users to make use of the objects in the schema, additional privileges might need to be granted, as appropriate for the object.

A user can also be allowed to create objects in someone else’s schema. To allow that, the CREATE privilege on the schema needs to be granted. Note that by default, everyone has CREATE and USAGE privileges on the schema public. This allows all users that are able to connect to a given database to create objects in its public schema. Some usage patterns call for revoking that privilege:

\[ \text{REVOKE CREATE ON SCHEMA public FROM PUBLIC;} \]

(The first “public” is the schema, the second “public” means “every user”. In the first sense it is an identifier, in the second sense it is a key word, hence the different capitalization; recall the guidelines from Section 4.1.1.)

### 5.8.5. The System Catalog Schema

In addition to public and user-created schemas, each database contains a pg_catalog schema, which contains the system tables and all the built-in data types, functions, and operators. pg_cat-
Data Definition

alog is always effectively part of the search path. If it is not named explicitly in the path then it is implicitly searched before searching the path's schemas. This ensures that built-in names will always be findable. However, you can explicitly place pg_catalog at the end of your search path if you prefer to have user-defined names override built-in names.

Since system table names begin with pg_, it is best to avoid such names to ensure that you won't suffer a conflict if some future version defines a system table named the same as your table. (With the default search path, an unqualified reference to your table name would then be resolved as the system table instead.) System tables will continue to follow the convention of having names beginning with pg_, so that they will not conflict with unqualified user-table names so long as users avoid the pg_ prefix.

5.8.6. Usage Patterns

Schemas can be used to organize your data in many ways. There are a few usage patterns easily supported by the default configuration, only one of which suffices when database users mistrust other database users:

• Constrain ordinary users to user-private schemas. To implement this, issue REVOKE CREATE ON SCHEMA public FROM PUBLIC, and create a schema for each user with the same name as that user. If affected users had logged in before this, consider auditing the public schema for objects named like objects in schema pg_catalog. Recall that the default search path starts with $user, which resolves to the user name. Therefore, if each user has a separate schema, they access their own schemas by default.

• Remove the public schema from each user's default search path using ALTER ROLE user SET search_path = "$user". Everyone retains the ability to create objects in the public schema, but only qualified names will choose those objects. While qualified table references are fine, calls to functions in the public schema will be unsafe or unreliable. Also, a user holding the CREATEROLE privilege can undo this setting and issue arbitrary queries under the identity of users relying on the setting. If you create functions or extensions in the public schema or grant CREATEROLE to users not warranting this almost-superuser ability, use the first pattern instead.

• Remove the public schema from search_path in postgresql.conf. The ensuing user experience matches the previous pattern. In addition to that pattern's implications for functions and CREATEROLE, this trusts database owners like CREATEROLE. If you create functions or extensions in the public schema or assign the CREATEROLE privilege, CREATEDB privilege or individual database ownership to users not warranting almost-superuser access, use the first pattern instead.

• Keep the default. All users access the public schema implicitly. This simulates the situation where schemas are not available at all, giving a smooth transition from the non-schema-aware world. However, any user can issue arbitrary queries under the identity of any user not electing to protect itself individually. This pattern is acceptable only when the database has a single user or a few mutually-trusting users.

For any pattern, to install shared applications (tables to be used by everyone, additional functions provided by third parties, etc.), put them into separate schemas. Remember to grant appropriate privileges to allow the other users to access them. Users can then refer to these additional objects by qualifying the names with a schema name, or they can put the additional schemas into their search path, as they choose.

5.8.7. Portability

In the SQL standard, the notion of objects in the same schema being owned by different users does not exist. Moreover, some implementations do not allow you to create schemas that have a different name than their owner. In fact, the concepts of schema and user are nearly equivalent in a database system that implements only the basic schema support specified in the standard. Therefore, many users consider qualified names to really consist of user_name.table_name. This is how PostgreSQL will effectively behave if you create a per-user schema for every user.
Also, there is no concept of a public schema in the SQL standard. For maximum conformance to the standard, you should not use the public schema.

Of course, some SQL database systems might not implement schemas at all, or provide namespace support by allowing (possibly limited) cross-database access. If you need to work with those systems, then maximum portability would be achieved by not using schemas at all.

## 5.9. Inheritance

PostgreSQL implements table inheritance, which can be a useful tool for database designers. (SQL:1999 and later define a type inheritance feature, which differs in many respects from the features described here.)

Let's start with an example: suppose we are trying to build a data model for cities. Each state has many cities, but only one capital. We want to be able to quickly retrieve the capital city for any particular state. This can be done by creating two tables, one for state capitals and one for cities that are not capitals. However, what happens when we want to ask for data about a city, regardless of whether it is a capital or not? The inheritance feature can help to resolve this problem. We define the capitals table so that it inherits from cities:

```sql
CREATE TABLE cities (  
  name text,  
  population float,  
  altitude int -- in feet
);

CREATE TABLE capitals (  
  state char(2)  
) INHERITS (cities);
```

In this case, the capitals table inherits all the columns of its parent table, cities. State capitals also have an extra column, state, that shows their state.

In PostgreSQL, a table can inherit from zero or more other tables, and a query can reference either all rows of a table or all rows of a table plus all of its descendant tables. The latter behavior is the default. For example, the following query finds the names of all cities, including state capitals, that are located at an altitude over 500 feet:

```sql
SELECT name, altitude  
FROM cities  
WHERE altitude > 500;
```

Given the sample data from the PostgreSQL tutorial (see Section 2.1), this returns:

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

On the other hand, the following query finds all the cities that are not state capitals and are situated at an altitude over 500 feet:

```sql
SELECT name, altitude  
FROM ONLY cities
```
WHERE altitude > 500;

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
</tbody>
</table>

Here the ONLY keyword indicates that the query should apply only to cities, and not any tables below cities in the inheritance hierarchy. Many of the commands that we have already discussed — SELECT, UPDATE and DELETE — support the ONLY keyword.

You can also write the table name with a trailing * to explicitly specify that descendant tables are included:

```sql
SELECT name, altitude
FROM cities*
WHERE altitude > 500;
```

Writing * is not necessary, since this behavior is always the default. However, this syntax is still supported for compatibility with older releases where the default could be changed.

In some cases you might wish to know which table a particular row originated from. There is a system column called tableoid in each table which can tell you the originating table:

```sql
SELECT c.tableoid, c.name, c.altitude
FROM cities c
WHERE c.altitude > 500;
```

which returns:

<table>
<thead>
<tr>
<th>tableoid</th>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>139793</td>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>139793</td>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>139798</td>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

(If you try to reproduce this example, you will probably get different numeric OIDs.) By doing a join with pg_class you can see the actual table names:

```sql
SELECT p.relname, c.name, c.altitude
FROM cities c, pg_class p
WHERE c.altitude > 500 AND c.tableoid = p.oid;
```

which returns:

<table>
<thead>
<tr>
<th>relname</th>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>cities</td>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>cities</td>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>capitals</td>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

Another way to get the same effect is to use the regclass alias type, which will print the table OID symbolically:
Data Definition

```
SELECT c.tableoid::regclass, c.name, c.altitude
FROM cities c
WHERE c.altitude > 500;
```

Inheritance does not automatically propagate data from INSERT or COPY commands to other tables in the inheritance hierarchy. In our example, the following INSERT statement will fail:

```
INSERT INTO cities (name, population, altitude, state)
VALUES ('Albany', NULL, NULL, 'NY');
```

We might hope that the data would somehow be routed to the capitals table, but this does not happen: INSERT always inserts into exactly the table specified. In some cases it is possible to redirect the insertion using a rule (see Chapter 41). However that does not help for the above case because the cities table does not contain the column state, and so the command will be rejected before the rule can be applied.

All check constraints and not-null constraints on a parent table are automatically inherited by its children, unless explicitly specified otherwise with NO INHERIT clauses. Other types of constraints (unique, primary key, and foreign key constraints) are not inherited.

A table can inherit from more than one parent table, in which case it has the union of the columns defined by the parent tables. Any columns declared in the child table's definition are added to these. If the same column name appears in multiple parent tables, or in both a parent table and the child's definition, then these columns are “merged” so that there is only one such column in the child table. To be merged, columns must have the same data types, else an error is raised. Inheritable check constraints and not-null constraints are merged in a similar fashion. Thus, for example, a merged column will be marked not-null if any one of the column definitions it came from is marked not-null. Check constraints are merged if they have the same name, and the merge will fail if their conditions are different.

Table inheritance is typically established when the child table is created, using the INHERITS clause of the CREATE TABLE statement. Alternatively, a table which is already defined in a compatible way can have a new parent relationship added, using the INHERIT variant of ALTER TABLE. To do this the new child table must already include columns with the same names and types as the columns of the parent. It must also include check constraints with the same names and check expressions as those of the parent. Similarly an inheritance link can be removed from a child using the NO INHERIT variant of ALTER TABLE. Dynamically adding and removing inheritance links like this can be useful when the inheritance relationship is being used for table partitioning (see Section 5.10).

One convenient way to create a compatible table that will later be made a new child is to use the LIKE clause in CREATE TABLE. This creates a new table with the same columns as the source table. If there are any CHECK constraints defined on the source table, the INCLUDING CONSTRAINTS option to LIKE should be specified, as the new child must have constraints matching the parent to be considered compatible.

A parent table cannot be dropped while any of its children remain. Neither can columns or check constraints of child tables be dropped or altered if they are inherited from any parent tables. If you wish to remove a table and all of its descendants, one easy way is to drop the parent table with the CASCADE option (see Section 5.13).

ALTER TABLE will propagate any changes in column data definitions and check constraints down the inheritance hierarchy. Again, dropping columns that are depended on by other tables is only possible when using the CASCADE option. ALTER TABLE follows the same rules for duplicate column merging and rejection that apply during CREATE TABLE.

Inherited queries perform access permission checks on the parent table only. Thus, for example, granting UPDATE permission on the cities table implies permission to update rows in the capitals table as well, when they are accessed through cities. This preserves the appearance that the data is (also) in the parent table. But the capitals table could not be updated directly without an addi-
tional grant. In a similar way, the parent table's row security policies (see Section 5.7) are applied to rows coming from child tables during an inherited query. A child table's policies, if any, are applied only when it is the table explicitly named in the query; and in that case, any policies attached to its parent(s) are ignored.

Foreign tables (see Section 5.11) can also be part of inheritance hierarchies, either as parent or child tables, just as regular tables can be. If a foreign table is part of an inheritance hierarchy then any operations not supported by the foreign table are not supported on the whole hierarchy either.

5.9.1. Caveats

Note that not all SQL commands are able to work on inheritance hierarchies. Commands that are used for data querying, data modification, or schema modification (e.g., `SELECT`, `UPDATE`, `DELETE`, most variants of `ALTER TABLE`, but not `INSERT` or `ALTER TABLE ... RENAME`) typically default to including child tables and support the `ONLY` notation to exclude them. Commands that do database maintenance and tuning (e.g., `REINDEX`, `VACUUM`) typically only work on individual, physical tables and do not support recursing over inheritance hierarchies. The respective behavior of each individual command is documented in its reference page (SQL Commands).

A serious limitation of the inheritance feature is that indexes (including unique constraints) and foreign key constraints only apply to single tables, not to their inheritance children. This is true on both the referencing and referenced sides of a foreign key constraint. Thus, in the terms of the above example:

- If we declared `cities.name` to be `UNIQUE` or a `PRIMARY KEY`, this would not stop the `capitals` table from having rows with names duplicating rows in `cities`. And those duplicate rows would by default show up in queries from `cities`. In fact, by default `capitals` would have no unique constraint at all, and so could contain multiple rows with the same name. You could add a unique constraint to `capitals`, but this would not prevent duplication compared to `cities`.

- Similarly, if we were to specify that `cities.name` REFERENCES some other table, this constraint would not automatically propagate to `capitals`. In this case you could work around it by manually adding the same REFERENCES constraint to `capitals`.

- Specifying that another table's column REFERENCES `cities(name)` would allow the other table to contain city names, but not capital names. There is no good workaround for this case.

These deficiencies will probably be fixed in some future release, but in the meantime considerable care is needed in deciding whether inheritance is useful for your application.

5.10. Table Partitioning

PostgreSQL supports basic table partitioning. This section describes why and how to implement partitioning as part of your database design.

5.10.1. Overview

Partitioning refers to splitting what is logically one large table into smaller physical pieces. Partitioning can provide several benefits:

- Query performance can be improved dramatically in certain situations, particularly when most of the heavily accessed rows of the table are in a single partition or a small number of partitions. The partitioning substitutes for leading columns of indexes, reducing index size and making it more likely that the heavily-used parts of the indexes fit in memory.

- When queries or updates access a large percentage of a single partition, performance can be improved by taking advantage of sequential scan of that partition instead of using an index and random access reads scattered across the whole table.
• Bulk loads and deletes can be accomplished by adding or removing partitions, if that requirement is planned into the partitioning design. Doing ALTER TABLE DETACH PARTITION or dropping an individual partition using DROP TABLE is far faster than a bulk operation. These commands also entirely avoid the VACUUM overhead caused by a bulk DELETE.

• Seldom-used data can be migrated to cheaper and slower storage media.

The benefits will normally be worthwhile only when a table would otherwise be very large. The exact point at which a table will benefit from partitioning depends on the application, although a rule of thumb is that the size of the table should exceed the physical memory of the database server.

PostgreSQL offers built-in support for the following forms of partitioning:

Range Partitioning

The table is partitioned into “ranges” defined by a key column or set of columns, with no overlap between the ranges of values assigned to different partitions. For example, one might partition by date ranges, or by ranges of identifiers for particular business objects.

List Partitioning

The table is partitioned by explicitly listing which key values appear in each partition.

Hash Partitioning

The table is partitioned by specifying a modulus and a remainder for each partition. Each partition will hold the rows for which the hash value of the partition key divided by the specified modulus will produce the specified remainder.

If your application needs to use other forms of partitioning not listed above, alternative methods such as inheritance and UNION ALL views can be used instead. Such methods offer flexibility but do not have some of the performance benefits of built-in declarative partitioning.

5.10.2. Declarative Partitioning

PostgreSQL offers a way to specify how to divide a table into pieces called partitions. The table that is divided is referred to as a partitioned table. The specification consists of the partitioning method and a list of columns or expressions to be used as the partition key.

All rows inserted into a partitioned table will be routed to one of the partitions based on the value of the partition key. Each partition has a subset of the data defined by its partition bounds. The currently supported partitioning methods are range, list, and hash.

Partitions may themselves be defined as partitioned tables, using what is called sub-partitioning. Partitions may have their own indexes, constraints and default values, distinct from those of other partitions. See CREATE TABLE for more details on creating partitioned tables and partitions.

It is not possible to turn a regular table into a partitioned table or vice versa. However, it is possible to add a regular or partitioned table containing data as a partition of a partitioned table, or remove a partition from a partitioned table turning it into a standalone table; see ALTER TABLE to learn more about the ATTACH PARTITION and DETACH PARTITION sub-commands.

Individual partitions are linked to the partitioned table with inheritance behind-the-scenes; however, it is not possible to use some of the generic features of inheritance (discussed below) with declaratively partitioned tables or their partitions. For example, a partition cannot have any parents other than the partitioned table it is a partition of, nor can a regular table inherit from a partitioned table making the latter its parent. That means partitioned tables and their partitions do not participate in inheritance with regular tables. Since a partition hierarchy consisting of the partitioned table and its partitions is
still an inheritance hierarchy, all the normal rules of inheritance apply as described in Section 5.9 with some exceptions, most notably:

- Both CHECK and NOT NULL constraints of a partitioned table are always inherited by all its partitions. CHECK constraints that are marked NO INHERIT are not allowed to be created on partitioned tables.

- Using ONLY to add or drop a constraint on only the partitioned table is supported as long as there are no partitions. Once partitions exist, using ONLY will result in an error as adding or dropping constraints on only the partitioned table, when partitions exist, is not supported. Instead, constraints on the partitions themselves can be added and (if they are not present in the parent table) dropped.

- As a partitioned table does not have any data directly, attempts to use TRUNCATE ONLY on a partitioned table will always return an error.

- Partitions cannot have columns that are not present in the parent. It is not possible to specify columns when creating partitions with CREATE TABLE, nor is it possible to add columns to partitions after-the-fact using ALTER TABLE. Tables may be added as a partition with ALTER TABLE ... ATTACH PARTITION only if their columns exactly match the parent, including any oid column.

- You cannot drop the NOT NULL constraint on a partition’s column if the constraint is present in the parent table.

Partitions can also be foreign tables, although they have some limitations that normal tables do not; see CREATE FOREIGN TABLE for more information.

Updating the partition key of a row might cause it to be moved into a different partition where this row satisfies the partition bounds.

### 5.10.2.1. Example

Suppose we are constructing a database for a large ice cream company. The company measures peak temperatures every day as well as ice cream sales in each region. Conceptually, we want a table like:

```sql
CREATE TABLE measurement (
    city_id int not null,
    logdate date not null,
    peaktemp int,
    unitsales int
);
```

We know that most queries will access just the last week’s, month’s or quarter’s data, since the main use of this table will be to prepare online reports for management. To reduce the amount of old data that needs to be stored, we decide to only keep the most recent 3 years worth of data. At the beginning of each month we will remove the oldest month’s data. In this situation we can use partitioning to help us meet all of our different requirements for the measurements table.

To use declarative partitioning in this case, use the following steps:

1. Create `measurement` table as a partitioned table by specifying the PARTITION BY clause, which includes the partitioning method (RANGE in this case) and the list of column(s) to use as the partition key.

```sql
CREATE TABLE measurement (  
    city_id int not null,
    logdate date not null,
    peaktemp int,
);
```
Data Definition

unitsales       | int |
) PARTITION BY RANGE (logdate);

You may decide to use multiple columns in the partition key for range partitioning, if desired. Of course, this will often result in a larger number of partitions, each of which is individually smaller. On the other hand, using fewer columns may lead to a coarser-grained partitioning criteria with smaller number of partitions. A query accessing the partitioned table will have to scan fewer partitions if the conditions involve some or all of these columns. For example, consider a table range partitioned using columns lastname and firstname (in that order) as the partition key.

2. Create partitions. Each partition’s definition must specify the bounds that correspond to the partitioning method and partition key of the parent. Note that specifying bounds such that the new partition’s values will overlap with those in one or more existing partitions will cause an error. Inserting data into the parent table that does not map to one of the existing partitions will cause an error; an appropriate partition must be added manually.

Partitions thus created are in every way normal PostgreSQL tables (or, possibly, foreign tables). It is possible to specify a tablespace and storage parameters for each partition separately.

It is not necessary to create table constraints describing partition boundary condition for partitions. Instead, partition constraints are generated implicitly from the partition bound specification whenever there is need to refer to them.

CREATE TABLE measurement_y2006m02 PARTITION OF measurement
   FOR VALUES FROM ('2006-02-01') TO ('2006-03-01');

CREATE TABLE measurement_y2006m03 PARTITION OF measurement
   FOR VALUES FROM ('2006-03-01') TO ('2006-04-01');

... 

CREATE TABLE measurement_y2007m11 PARTITION OF measurement
   FOR VALUES FROM ('2007-11-01') TO ('2007-12-01');

CREATE TABLE measurement_y2007m12 PARTITION OF measurement
   FOR VALUES FROM ('2007-12-01') TO ('2008-01-01')
   TABLESPACE fasttablespace;

CREATE TABLE measurement_y2008m01 PARTITION OF measurement
   FOR VALUES FROM ('2008-01-01') TO ('2008-02-01')
   WITH (parallel_workers = 4)
   TABLESPACE fasttablespace;

To implement sub-partitioning, specify the PARTITION BY clause in the commands used to create individual partitions, for example:

CREATE TABLE measurement_y2006m02 PARTITION OF measurement
   FOR VALUES FROM ('2006-02-01') TO ('2006-03-01')
   PARTITION BY RANGE (peaktemp);

After creating partitions of measurement_y2006m02, any data inserted into measurement that is mapped to measurement_y2006m02 (or data that is directly inserted into measurement_y2006m02, provided it satisfies its partition constraint) will be further redirected to one of its partitions based on the peaktemp column. The partition key specified may overlap with the parent’s partition key, although care should be taken when specifying the bounds of a sub-partition such that the set of data it accepts constitutes a subset of what the partition’s own bounds allows; the system does not try to check whether that’s really the case.

3. Create an index on the key column(s), as well as any other indexes you might want, on the partitioned table. (The key index is not strictly necessary, but in most scenarios it is helpful.)
automatically creates one index on each partition, and any partitions you create or attach later will also contain the index.

```
CREATE INDEX ON measurement (logdate);
```

4. Ensure that the `enable_partition_pruning` configuration parameter is not disabled in `postgresql.conf`. If it is, queries will not be optimized as desired.

In the above example we would be creating a new partition each month, so it might be wise to write a script that generates the required DDL automatically.

### 5.10.2.2. Partition Maintenance

Normally the set of partitions established when initially defining the table are not intended to remain static. It is common to want to remove old partitions of data and periodically add new partitions for new data. One of the most important advantages of partitioning is precisely that it allows this otherwise painful task to be executed nearly instantaneously by manipulating the partition structure, rather than physically moving large amounts of data around.

The simplest option for removing old data is to drop the partition that is no longer necessary:

```
DROP TABLE measurement_y2006m02;
```

This can very quickly delete millions of records because it doesn't have to individually delete every record. Note however that the above command requires taking an `ACCESS EXCLUSIVE` lock on the parent table.

Another option that is often preferable is to remove the partition from the partitioned table but retain access to it as a table in its own right:

```
ALTER TABLE measurement DETACH PARTITION measurement_y2006m02;
```

This allows further operations to be performed on the data before it is dropped. For example, this is often a useful time to back up the data using `COPY`, `pg_dump`, or similar tools. It might also be a useful time to aggregate data into smaller formats, perform other data manipulations, or run reports.

Similarly we can add a new partition to handle new data. We can create an empty partition in the partitioned table just as the original partitions were created above:

```
CREATE TABLE measurement_y2008m02 PARTITION OF measurement
    FOR VALUES FROM ('2008-02-01') TO ('2008-03-01')
    TABLESPACE fasttablespace;
```

As an alternative, it is sometimes more convenient to create the new table outside the partition structure, and make it a proper partition later. This allows the data to be loaded, checked, and transformed prior to it appearing in the partitioned table:

```
CREATE TABLE measurement_y2008m02
    (LIKE measurement INCLUDING DEFAULTS INCLUDING CONSTRAINTS)
    TABLESPACE fasttablespace;

ALTER TABLE measurement_y2008m02 ADD CONSTRAINT y2008m02
    CHECK ( logdate >= DATE '2008-02-01' AND logdate < DATE '2008-03-01' );

\copy measurement_y2008m02 from 'measurement_y2008m02'
```
-- possibly some other data preparation work

ALTER TABLE measurement ATTACH PARTITION measurement_y2008m02
    FOR VALUES FROM ('2008-02-01') TO ('2008-03-01');

Before running the ATTACH PARTITION command, it is recommended to create a CHECK constraint on the table to be attached describing the desired partition constraint. That way, the system will be able to skip the scan to validate the implicit partition constraint. Without such a constraint, the table will be scanned to validate the partition constraint while holding an ACCESS EXCLUSIVE lock on the parent table. One may then drop the constraint after ATTACH PARTITION is finished, because it is no longer necessary.

5.10.2.3. Limitations

The following limitations apply to partitioned tables:

- There is no way to create an exclusion constraint spanning all partitions; it is only possible to constrain each leaf partition individually.

- While primary keys are supported on partitioned tables, foreign keys referencing partitioned tables are not supported. (Foreign key references from a partitioned table to some other table are supported.)

- When an UPDATE causes a row to move from one partition to another, there is a chance that another concurrent UPDATE or DELETE misses this row. Suppose session 1 is performing an UPDATE on a partition key, and meanwhile a concurrent session 2 for which this row is visible performs an UPDATE or DELETE operation on this row. Session 2 can silently miss the row if the row is deleted from the partition due to session 1's activity. In such case, session 2's UPDATE or DELETE, being unaware of the row movement thinks that the row has just been deleted and concludes that there is nothing to be done for this row. In the usual case where the table is not partitioned, or where there is no row movement, session 2 would have identified the newly updated row and carried out the UPDATE/DELETE on this new row version.

- BEFORE ROW triggers, if necessary, must be defined on individual partitions, not the partitioned table.

- Mixing temporary and permanent relations in the same partition tree is not allowed. Hence, if the partitioned table is permanent, so must be its partitions and likewise if the partitioned table is temporary. When using temporary relations, all members of the partition tree have to be from the same session.

5.10.3. Implementation Using Inheritance

While the built-in declarative partitioning is suitable for most common use cases, there are some circumstances where a more flexible approach may be useful. Partitioning can be implemented using table inheritance, which allows for several features not supported by declarative partitioning, such as:

- For declarative partitioning, partitions must have exactly the same set of columns as the partitioned table, whereas with table inheritance, child tables may have extra columns not present in the parent.

- Table inheritance allows for multiple inheritance.

- Declarative partitioning only supports range, list and hash partitioning, whereas table inheritance allows data to be divided in a manner of the user's choosing. (Note, however, that if constraint exclusion is unable to prune child tables effectively, query performance might be poor.)

- Some operations require a stronger lock when using declarative partitioning than when using table inheritance. For example, adding or removing a partition to or from a partitioned table requires taking an ACCESS EXCLUSIVE lock on the parent table, whereas a SHARE UPDATE EXCLUSIVE lock is enough in the case of regular inheritance.
5.10.3.1. Example

We use the same measurement table we used above. To implement partitioning using inheritance, use the following steps:

1. Create the “master” table, from which all of the “child” tables will inherit. This table will contain no data. Do not define any check constraints on this table, unless you intend them to be applied equally to all child tables. There is no point in defining any indexes or unique constraints on it, either. For our example, the master table is the measurement table as originally defined.

2. Create several “child” tables that each inherit from the master table. Normally, these tables will not add any columns to the set inherited from the master. Just as with declarative partitioning, these tables are in every way normal PostgreSQL tables (or foreign tables).

   CREATE TABLE measurement_y2006m02 () INHERITS (measurement);
   CREATE TABLE measurement_y2006m03 () INHERITS (measurement);
   ...
   CREATE TABLE measurement_y2007m11 () INHERITS (measurement);
   CREATE TABLE measurement_y2007m12 () INHERITS (measurement);
   CREATE TABLE measurement_y2008m01 () INHERITS (measurement);

3. Add non-overlapping table constraints to the child tables to define the allowed key values in each. Typical examples would be:

   CHECK ( x = 1 )
   CHECK ( county IN ( 'Oxfordshire', 'Buckinghamshire', 'Warwickshire' ))
   CHECK ( outletID >= 100 AND outletID < 200 )

Ensure that the constraints guarantee that there is no overlap between the key values permitted in different child tables. A common mistake is to set up range constraints like:

   CHECK ( outletID BETWEEN 100 AND 200 )
   CHECK ( outletID BETWEEN 200 AND 300 )

This is wrong since it is not clear which child table the key value 200 belongs in.

It would be better to instead create child tables as follows:

   CREATE TABLE measurement_y2006m02 (  
       CHECK ( logdate >= DATE '2006-02-01' AND logdate < DATE '2006-03-01' )  
    ) INHERITS (measurement);
   CREATE TABLE measurement_y2006m03 (  
       CHECK ( logdate >= DATE '2006-03-01' AND logdate < DATE '2006-04-01' )  
    ) INHERITS (measurement);
   ...
   CREATE TABLE measurement_y2007m11 (  
       CHECK ( logdate >= DATE '2007-11-01' AND logdate < DATE '2007-12-01' )  
    ) INHERITS (measurement);
   CREATE TABLE measurement_y2007m12 (  

CREATE TABLE measurement_y2008m01 (  
    CHECK ( logdate >= DATE '2008-01-01' AND logdate < DATE '2008-02-01' )  
) INHERITS (measurement);

4. For each child table, create an index on the key column(s), as well as any other indexes you might want.

CREATE INDEX measurement_y2006m02_logdate ON measurement_y2006m02 (logdate);
CREATE INDEX measurement_y2006m03_logdate ON measurement_y2006m03 (logdate);
CREATE INDEX measurement_y2007m11_logdate ON measurement_y2007m11 (logdate);
CREATE INDEX measurement_y2007m12_logdate ON measurement_y2007m12 (logdate);
CREATE INDEX measurement_y2008m01_logdate ON measurement_y2008m01 (logdate);

5. We want our application to be able to say INSERT INTO measurement ... and have the data be redirected into the appropriate child table. We can arrange that by attaching a suitable trigger function to the master table. If data will be added only to the latest child, we can use a very simple trigger function:

CREATE OR REPLACE FUNCTION measurement_insert_trigger()  
RETURNS TRIGGER AS $$  
BEGIN  
    INSERT INTO measurement_y2008m01 VALUES (NEW.*);
    RETURN NULL;
END;  
$$  
LANGUAGE plpgsql;

After creating the function, we create a trigger which calls the trigger function:

CREATE TRIGGER insert_measurement_trigger  
BEFORE INSERT ON measurement  
FOR EACH ROW EXECUTE FUNCTION measurement_insert_trigger();

We must redefine the trigger function each month so that it always points to the current child table. The trigger definition does not need to be updated, however.

We might want to insert data and have the server automatically locate the child table into which the row should be added. We could do this with a more complex trigger function, for example:

CREATE OR REPLACE FUNCTION measurement_insert_trigger()  
RETURNS TRIGGER AS $$  
BEGIN  
    IF ( NEW.logdate >= DATE '2006-02-01' AND  
        NEW.logdate < DATE '2006-03-01' ) THEN  
        INSERT INTO measurement_y2006m02 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-03-01' AND  
        NEW.logdate < DATE '2006-04-01' ) THEN  
        INSERT INTO measurement_y2006m03 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-04-01' AND  
        NEW.logdate < DATE '2006-05-01' ) THEN  
        INSERT INTO measurement_y2006m04 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-05-01' AND  
        NEW.logdate < DATE '2006-06-01' ) THEN  
        INSERT INTO measurement_y2006m05 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-06-01' AND  
        NEW.logdate < DATE '2006-07-01' ) THEN  
        INSERT INTO measurement_y2006m06 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-07-01' AND  
        NEW.logdate < DATE '2006-08-01' ) THEN  
        INSERT INTO measurement_y2006m07 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-08-01' AND  
        NEW.logdate < DATE '2006-09-01' ) THEN  
        INSERT INTO measurement_y2006m08 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-09-01' AND  
        NEW.logdate < DATE '2006-10-01' ) THEN  
        INSERT INTO measurement_y2006m09 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-10-01' AND  
        NEW.logdate < DATE '2006-11-01' ) THEN  
        INSERT INTO measurement_y2006m10 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-11-01' AND  
        NEW.logdate < DATE '2006-12-01' ) THEN  
        INSERT INTO measurement_y2006m11 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-12-01' AND  
        NEW.logdate < DATE '2007-01-01' ) THEN  
        INSERT INTO measurement_y2007m01 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-01-01' AND  
        NEW.logdate < DATE '2007-02-01' ) THEN  
        INSERT INTO measurement_y2007m02 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-02-01' AND  
        NEW.logdate < DATE '2007-03-01' ) THEN  
        INSERT INTO measurement_y2007m03 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-03-01' AND  
        NEW.logdate < DATE '2007-04-01' ) THEN  
        INSERT INTO measurement_y2007m04 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-04-01' AND  
        NEW.logdate < DATE '2007-05-01' ) THEN  
        INSERT INTO measurement_y2007m05 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-05-01' AND  
        NEW.logdate < DATE '2007-06-01' ) THEN  
        INSERT INTO measurement_y2007m06 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-06-01' AND  
        NEW.logdate < DATE '2007-07-01' ) THEN  
        INSERT INTO measurement_y2007m07 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-07-01' AND  
        NEW.logdate < DATE '2007-08-01' ) THEN  
        INSERT INTO measurement_y2007m08 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-08-01' AND  
        NEW.logdate < DATE '2007-09-01' ) THEN  
        INSERT INTO measurement_y2007m09 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-09-01' AND  
        NEW.logdate < DATE '2007-10-01' ) THEN  
        INSERT INTO measurement_y2007m10 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-10-01' AND  
        NEW.logdate < DATE '2007-11-01' ) THEN  
        INSERT INTO measurement_y2007m11 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2007-11-01' AND  
        NEW.logdate < DATE '2007-12-01' ) THEN  
        INSERT INTO measurement_y2007m12 VALUES (NEW.*);
    ELSE  
        INSERT INTO measurement VALUES (NEW.*);
    END;  
END;  
$$  
LANGUAGE plpgsql;
CREATE RULE measurement_insert_y2006m02 AS
ON INSERT TO measurement WHERE
  ( logdate >= DATE '2006-02-01' AND logdate < DATE '2006-03-01' )
DO INSTEAD
  INSERT INTO measurement_y2006m02 VALUES (NEW.*);...

CREATE RULE measurement_insert_y2008m01 AS
ON INSERT TO measurement WHERE
  ( logdate >= DATE '2008-01-01' AND logdate < DATE '2008-02-01' )
DO INSTEAD
  INSERT INTO measurement_y2008m01 VALUES (NEW.*);
As we can see, a complex table hierarchy could require a substantial amount of DDL. In the above example we would be creating a new child table each month, so it might be wise to write a script that generates the required DDL automatically.

5.10.3.2. Maintenance for Inheritance Partitioning

To remove old data quickly, simply drop the child table that is no longer necessary:

```
DROP TABLE measurement_y2006m02;
```

To remove the child table from the inheritance hierarchy table but retain access to it as a table in its own right:

```
ALTER TABLE measurement_y2006m02 NO INHERIT measurement;
```

To add a new child table to handle new data, create an empty child table just as the original children were created above:

```
CREATE TABLE measurement_y2008m02 (  
  CHECK ( logdate >= DATE '2008-02-01' AND logdate < DATE '2008-03-01' )  
) INHERITS (measurement);
```

Alternatively, one may want to create and populate the new child table before adding it to the table hierarchy. This could allow data to be loaded, checked, and transformed before being made visible to queries on the parent table.

```
CREATE TABLE measurement_y2008m02  
  (LIKE measurement INCLUDING DEFAULTS INCLUDING CONSTRAINTS);  
ALTER TABLE measurement_y2008m02 ADD CONSTRAINT y2008m02  
  CHECK ( logdate >= DATE '2008-02-01' AND logdate < DATE '2008-03-01' );  
\copy measurement_y2008m02 from 'measurement_y2008m02'  
-- possibly some other data preparation work  
ALTER TABLE measurement_y2008m02 INHERIT measurement;
```

5.10.3.3. Caveats

The following caveats apply to partitioning implemented using inheritance:

- There is no automatic way to verify that all of the \texttt{CHECK} constraints are mutually exclusive. It is safer to create code that generates child tables and creates and/or modifies associated objects than to write each by hand.

- The schemes shown here assume that the values of a row's key column(s) never change, or at least do not change enough to require it to move to another partition. An \texttt{UPDATE} that attempts to do that will fail because of the \texttt{CHECK} constraints. If you need to handle such cases, you can put suitable update triggers on the child tables, but it makes management of the structure much more complicated.

- If you are using manual \texttt{VACUUM} or \texttt{ANALYZE} commands, don't forget that you need to run them on each child table individually. A command like:

```
ANALYZE measurement;
```

will only process the master table.
• INSERT statements with ON CONFLICT clauses are unlikely to work as expected, as the ON CONFLICT action is only taken in case of unique violations on the specified target relation, not its child relations.

• Triggers or rules will be needed to route rows to the desired child table, unless the application is explicitly aware of the partitioning scheme. Triggers may be complicated to write, and will be much slower than the tuple routing performed internally by declarative partitioning.

5.10.4. Partition Pruning

Partition pruning is a query optimization technique that improves performance for declaratively partitioned tables. As an example:

```
SET enable_partition_pruning = on;          -- the default
SELECT count(*) FROM measurement WHERE logdate >= DATE '2008-01-01';
```

Without partition pruning, the above query would scan each of the partitions of the measurement table. With partition pruning enabled, the planner will examine the definition of each partition and prove that the partition need not be scanned because it could not contain any rows meeting the query's WHERE clause. When the planner can prove this, it excludes (prunes) the partition from the query plan.

By using the EXPLAIN command and the enable_partition_pruning configuration parameter, it's possible to show the difference between a plan for which partitions have been pruned and one for which they have not. A typical unoptimized plan for this type of table setup is:

```
SET enable_partition_pruning = off;
EXPLAIN SELECT count(*) FROM measurement WHERE logdate >= DATE '2008-01-01';
```

```
QUERY PLAN
-----------------------------------------------------------------------------------
Aggregate  (cost=188.76..188.77 rows=1 width=8)
  ->  Append  (cost=0.00..181.05 rows=3085 width=0)
    ->  Seq Scan on measurement_y2006m02  (cost=0.00..33.12 rows=617 width=0)
          Filter: (logdate >= '2008-01-01'::date)
    ->  Seq Scan on measurement_y2006m03  (cost=0.00..33.12 rows=617 width=0)
          Filter: (logdate >= '2008-01-01'::date)
          ...  
    ->  Seq Scan on measurement_y2007m11  (cost=0.00..33.12 rows=617 width=0)
          Filter: (logdate >= '2008-01-01'::date)
    ->  Seq Scan on measurement_y2007m12  (cost=0.00..33.12 rows=617 width=0)
          Filter: (logdate >= '2008-01-01'::date)
    ->  Seq Scan on measurement_y2008m01  (cost=0.00..33.12 rows=617 width=0)
          Filter: (logdate >= '2008-01-01'::date)
```

Some or all of the partitions might use index scans instead of full-table sequential scans, but the point here is that there is no need to scan the older partitions at all to answer this query. When we enable partition pruning, we get a significantly cheaper plan that will deliver the same answer:

```
SET enable_partition_pruning = on;
EXPLAIN SELECT count(*) FROM measurement WHERE logdate >= DATE '2008-01-01';
```
Partition pruning can be performed not only during the planning of a given query, but also during its execution. This is useful as it can allow more partitions to be pruned when clauses contain expressions whose values are not known at query planning time; for example, parameters defined in a PREPARE statement, using a value obtained from a subquery or using a parameterized value on the inner side of a nested loop join. Partition pruning during execution can be performed at any of the following times:

- During initialization of the query plan. Partition pruning can be performed here for parameter values which are known during the initialization phase of execution. Partitions which are pruned during this stage will not show up in the query’s EXPLAIN or EXPLAIN ANALYZE. It is possible to determine the number of partitions which were removed during this phase by observing the “Subplans Removed” property in the EXPLAIN output.

- During actual execution of the query plan. Partition pruning may also be performed here to remove partitions using values which are only known during actual query execution. This includes values from subqueries and values from execution-time parameters such as those from parameterized nested loop joins. Since the value of these parameters may change many times during the execution of the query, partition pruning is performed whenever one of the execution parameters being used by partition pruning changes. Determining if partitions were pruned during this phase requires careful inspection of the nloops property in the EXPLAIN ANALYZE output.

Partition pruning can be disabled using the enable_partition_pruning setting.

Note

Currently, pruning of partitions during the planning of an UPDATE or DELETE command is implemented using the constraint exclusion method (however, it is controlled by the enable_partition_pruning rather than constraint_exclusion) — see the following section for details and caveats that apply.

Also, execution-time partition pruning currently only occurs for the Append node type, not MergeAppend.

Both of these behaviors are likely to be changed in a future release of PostgreSQL.

5.10.5. Partitioning and Constraint Exclusion

Constraint exclusion is a query optimization technique similar to partition pruning. While it is primarily used for partitioning implemented using the legacy inheritance method, it can be used for other purposes, including with declarative partitioning.

Constraint exclusion works in a very similar way to partition pruning, except that it uses each table’s CHECK constraints — which gives it its name — whereas partition pruning uses the table’s partition bounds, which exist only in the case of declarative partitioning. Another difference is that constraint exclusion is only applied at plan time; there is no attempt to remove partitions at execution time.
The fact that constraint exclusion uses \texttt{CHECK} constraints, which makes it slow compared to partition pruning, can sometimes be used as an advantage: because constraints can be defined even on declaratively-partitioned tables, in addition to their internal partition bounds, constraint exclusion may be able to elide additional partitions from the query plan.

The default (and recommended) setting of constraint\_exclusion is neither \texttt{on} nor \texttt{off}, but an intermediate setting called \texttt{partition}, which causes the technique to be applied only to queries that are likely to be working on inheritance partitioned tables. The \texttt{on} setting causes the planner to examine \texttt{CHECK} constraints in all queries, even simple ones that are unlikely to benefit.

The following caveats apply to constraint exclusion:

- Constraint exclusion is only applied during query planning; unlike partition pruning, it cannot be applied during query execution.

- Constraint exclusion only works when the query's \texttt{WHERE} clause contains constants (or externally supplied parameters). For example, a comparison against a non-immutable function such as \texttt{CURRENT\_TIMESTAMP} cannot be optimized, since the planner cannot know which child table the function's value might fall into at run time.

- Keep the partitioning constraints simple, else the planner may not be able to prove that child tables might not need to be visited. Use simple equality conditions for list partitioning, or simple range tests for range partitioning, as illustrated in the preceding examples. A good rule of thumb is that partitioning constraints should contain only comparisons of the partitioning column(s) to constants using B-tree-indexable operators, because only B-tree-indexable column(s) are allowed in the partition key.

- All constraints on all children of the parent table are examined during constraint exclusion, so large numbers of children are likely to increase query planning time considerably. So the legacy inheritance based partitioning will work well with up to perhaps a hundred child tables; don't try to use many thousands of children.

### 5.11. Foreign Data

PostgreSQL implements portions of the SQL/MED specification, allowing you to access data that resides outside PostgreSQL using regular SQL queries. Such data is referred to as foreign data. (Note that this usage is not to be confused with foreign keys, which are a type of constraint within the database.)

Foreign data is accessed with help from a foreign data wrapper. A foreign data wrapper is a library that can communicate with an external data source, hiding the details of connecting to the data source and obtaining data from it. There are some foreign data wrappers available as contrib modules; see Appendix F. Other kinds of foreign data wrappers might be found as third party products. If none of the existing foreign data wrappers suit your needs, you can write your own; see Chapter 57.

To access foreign data, you need to create a foreign server object, which defines how to connect to a particular external data source according to the set of options used by its supporting foreign data wrapper. Then you need to create one or more foreign tables, which define the structure of the remote data. A foreign table can be used in queries just like a normal table, but a foreign table has no storage in the PostgreSQL server. Whenever it is used, PostgreSQL asks the foreign data wrapper to fetch data from the external source, or transmit data to the external source in the case of update commands.

Accessing remote data may require authenticating to the external data source. This information can be provided by a user mapping, which can provide additional data such as user names and passwords based on the current PostgreSQL role.

For additional information, see CREATE FOREIGN DATA WRAPPER, CREATE SERVER, CREATE USER MAPPING, CREATE FOREIGN TABLE, and IMPORT FOREIGN SCHEMA.
5.12. Other Database Objects

Tables are the central objects in a relational database structure, because they hold your data. But they are not the only objects that exist in a database. Many other kinds of objects can be created to make the use and management of the data more efficient or convenient. They are not discussed in this chapter, but we give you a list here so that you are aware of what is possible:

- Views
- Functions, procedures, and operators
- Data types and domains
- Triggers and rewrite rules

Detailed information on these topics appears in Part V.

5.13. Dependency Tracking

When you create complex database structures involving many tables with foreign key constraints, views, triggers, functions, etc. you implicitly create a net of dependencies between the objects. For instance, a table with a foreign key constraint depends on the table it references.

To ensure the integrity of the entire database structure, PostgreSQL makes sure that you cannot drop objects that other objects still depend on. For example, attempting to drop the products table we considered in Section 5.3.5, with the orders table depending on it, would result in an error message like this:

```
DROP TABLE products;
ERROR:  cannot drop table products because other objects depend on it
DETAIL:  constraint orders_product_no_fkey on table orders depends on table products
HINT:  Use DROP ... CASCADE to drop the dependent objects too.
```

The error message contains a useful hint: if you do not want to bother deleting all the dependent objects individually, you can run:

```
DROP TABLE products CASCADE;
```

and all the dependent objects will be removed, as will any objects that depend on them, recursively. In this case, it doesn’t remove the orders table, it only removes the foreign key constraint. It stops there because nothing depends on the foreign key constraint. (If you want to check what `DROP ... CASCADE` will do, run `DROP` without `CASCADE` and read the `DETAIL` output.)

Almost all `DROP` commands in PostgreSQL support specifying `CASCADE`. Of course, the nature of the possible dependencies varies with the type of the object. You can also write `RESTRICT` instead of `CASCADE` to get the default behavior, which is to prevent dropping objects that any other objects depend on.

```
Note
According to the SQL standard, specifying either RESTRICT or CASCADE is required in a DROP command. No database system actually enforces that rule, but whether the default behavior is RESTRICT or CASCADE varies across systems.
```
If a DROP command lists multiple objects, CASCADE is only required when there are dependencies outside the specified group. For example, when saying DROP TABLE tab1, tab2 the existence of a foreign key referencing tab1 from tab2 would not mean that CASCADE is needed to succeed.

For user-defined functions, PostgreSQL tracks dependencies associated with a function's externally-visible properties, such as its argument and result types, but not dependencies that could only be known by examining the function body. As an example, consider this situation:

```sql
CREATE TYPE rainbow AS ENUM ('red', 'orange', 'yellow',
  'green', 'blue', 'purple');

CREATE TABLE my_colors (color rainbow, note text);

CREATE FUNCTION get_color_note (rainbow) RETURNS text AS
  'SELECT note FROM my_colors WHERE color = $1'
  LANGUAGE SQL;
```

(See Section 38.5 for an explanation of SQL-language functions.) PostgreSQL will be aware that the `get_color_note` function depends on the `rainbow` type: dropping the type would force dropping the function, because its argument type would no longer be defined. But PostgreSQL will not consider `get_color_note` to depend on the `my_colors` table, and so will not drop the function if the table is dropped. While there are disadvantages to this approach, there are also benefits. The function is still valid in some sense if the table is missing, though executing it would cause an error; creating a new table of the same name would allow the function to work again.
Chapter 6. Data Manipulation

The previous chapter discussed how to create tables and other structures to hold your data. Now it is time to fill the tables with data. This chapter covers how to insert, update, and delete table data. The chapter after this will finally explain how to extract your long-lost data from the database.

6.1. Inserting Data

When a table is created, it contains no data. The first thing to do before a database can be of much use is to insert data. Data is conceptually inserted one row at a time. Of course you can also insert more than one row, but there is no way to insert less than one row. Even if you know only some column values, a complete row must be created.

To create a new row, use the INSERT command. The command requires the table name and column values. For example, consider the products table from Chapter 5:

```
CREATE TABLE products (
   product_no integer,
   name text,
   price numeric
);
```

An example command to insert a row would be:

```
INSERT INTO products VALUES (1, 'Cheese', 9.99);
```

The data values are listed in the order in which the columns appear in the table, separated by commas. Usually, the data values will be literals (constants), but scalar expressions are also allowed.

The above syntax has the drawback that you need to know the order of the columns in the table. To avoid this you can also list the columns explicitly. For example, both of the following commands have the same effect as the one above:

```
INSERT INTO products (product_no, name, price) VALUES (1, 'Cheese', 9.99);
INSERT INTO products (name, price, product_no) VALUES ('Cheese', 9.99, 1);
```

Many users consider it good practice to always list the column names.

If you don't have values for all the columns, you can omit some of them. In that case, the columns will be filled with their default values. For example:

```
INSERT INTO products (product_no, name) VALUES (1, 'Cheese');
INSERT INTO products VALUES (1, 'Cheese');
```

The second form is a PostgreSQL extension. It fills the columns from the left with as many values as are given, and the rest will be defaulted.

For clarity, you can also request default values explicitly, for individual columns or for the entire row:

```
INSERT INTO products (product_no, name, price) VALUES (1, 'Cheese', DEFAULT);
```
INSERT INTO products DEFAULT VALUES;

You can insert multiple rows in a single command:

```sql
INSERT INTO products (product_no, name, price) VALUES
    (1, 'Cheese', 9.99),
    (2, 'Bread', 1.99),
    (3, 'Milk', 2.99);
```

It is also possible to insert the result of a query (which might be no rows, one row, or many rows):

```sql
INSERT INTO products (product_no, name, price)
    SELECT product_no, name, price FROM new_products
    WHERE release_date = 'today';
```

This provides the full power of the SQL query mechanism (Chapter 7) for computing the rows to be inserted.

**Tip**

When inserting a lot of data at the same time, consider using the COPY command. It is not as flexible as the INSERT command, but is more efficient. Refer to Section 14.4 for more information on improving bulk loading performance.

### 6.2. Updating Data

The modification of data that is already in the database is referred to as updating. You can update individual rows, all the rows in a table, or a subset of all rows. Each column can be updated separately; the other columns are not affected.

To update existing rows, use the UPDATE command. This requires three pieces of information:

1. The name of the table and column to update
2. The new value of the column
3. Which row(s) to update

Recall from Chapter 5 that SQL does not, in general, provide a unique identifier for rows. Therefore it is not always possible to directly specify which row to update. Instead, you specify which conditions a row must meet in order to be updated. Only if you have a primary key in the table (independent of whether you declared it or not) can you reliably address individual rows by choosing a condition that matches the primary key. Graphical database access tools rely on this fact to allow you to update rows individually.

For example, this command updates all products that have a price of 5 to have a price of 10:

```sql
UPDATE products SET price = 10 WHERE price = 5;
```

This might cause zero, one, or many rows to be updated. It is not an error to attempt an update that does not match any rows.

Let’s look at that command in detail. First is the key word `UPDATE` followed by the table name. As usual, the table name can be schema-qualified, otherwise it is looked up in the path. Next is the key word `SET` followed by the column name, an equal sign, and the new column value. The new column value can be any scalar expression, not just a constant. For example, if you want to raise the price of all products by 10% you could use:
UPDATE products SET price = price * 1.10;

As you see, the expression for the new value can refer to the existing value(s) in the row. We also left out the WHERE clause. If it is omitted, it means that all rows in the table are updated. If it is present, only those rows that match the WHERE condition are updated. Note that the equals sign in the SET clause is an assignment while the one in the WHERE clause is a comparison, but this does not create any ambiguity. Of course, the WHERE condition does not have to be an equality test. Many other operators are available (see Chapter 9). But the expression needs to evaluate to a Boolean result.

You can update more than one column in an UPDATE command by listing more than one assignment in the SET clause. For example:

UPDATE mytable SET a = 5, b = 3, c = 1 WHERE a > 0;

6.3. Deleting Data

So far we have explained how to add data to tables and how to change data. What remains is to discuss how to remove data that is no longer needed. Just as adding data is only possible in whole rows, you can only remove entire rows from a table. In the previous section we explained that SQL does not provide a way to directly address individual rows. Therefore, removing rows can only be done by specifying conditions that the rows to be removed have to match. If you have a primary key in the table then you can specify the exact row. But you can also remove groups of rows matching a condition, or you can remove all rows in the table at once.

You use the DELETE command to remove rows; the syntax is very similar to the UPDATE command. For instance, to remove all rows from the products table that have a price of 10, use:

DELETE FROM products WHERE price = 10;

If you simply write:

DELETE FROM products;

then all rows in the table will be deleted! Caveat programmer.

6.4. Returning Data From Modified Rows

Sometimes it is useful to obtain data from modified rows while they are being manipulated. The INSERT, UPDATE, and DELETE commands all have an optional RETURNING clause that supports this. Use of RETURNING avoids performing an extra database query to collect the data, and is especially valuable when it would otherwise be difficult to identify the modified rows reliably.

The allowed contents of a RETURNING clause are the same as a SELECT command’s output list (see Section 7.3). It can contain column names of the command’s target table, or value expressions using those columns. A common shorthand is RETURNING *, which selects all columns of the target table in order.

In an INSERT, the data available to RETURNING is the row as it was inserted. This is not so useful in trivial inserts, since it would just repeat the data provided by the client. But it can be very handy when relying on computed default values. For example, when using a serial column to provide unique identifiers, RETURNING can return the ID assigned to a new row:

CREATE TABLE users (firstname text, lastname text, id serial primary key);
INSERT INTO users (firstname, lastname) VALUES ('Joe', 'Cool')
RETURNING id;

The RETURNING clause is also very useful with INSERT ... SELECT.

In an UPDATE, the data available to RETURNING is the new content of the modified row. For example:

UPDATE products SET price = price * 1.10
WHERE price <= 99.99
RETURNING name, price AS new_price;

In a DELETE, the data available to RETURNING is the content of the deleted row. For example:

DELETE FROM products
WHERE obsoletion_date = 'today'
RETURNING *;

If there are triggers (Chapter 39) on the target table, the data available to RETURNING is the row as modified by the triggers. Thus, inspecting columns computed by triggers is another common use-case for RETURNING.
Chapter 7. Queries

The previous chapters explained how to create tables, how to fill them with data, and how to manipulate that data. Now we finally discuss how to retrieve the data from the database.

7.1. Overview

The process of retrieving or the command to retrieve data from a database is called a query. In SQL the SELECT command is used to specify queries. The general syntax of the SELECT command is

```
[WITH with_queries] SELECT select_list FROM table_expression
[sort_specification]
```

The following sections describe the details of the select list, the table expression, and the sort specification. WITH queries are treated last since they are an advanced feature.

A simple kind of query has the form:

```
SELECT * FROM table1;
```

Assuming that there is a table called table1, this command would retrieve all rows and all user-defined columns from table1. (The method of retrieval depends on the client application. For example, the psql program will display an ASCII-art table on the screen, while client libraries will offer functions to extract individual values from the query result.) The select list specification * means all columns that the table expression happens to provide. A select list can also select a subset of the available columns or make calculations using the columns. For example, if table1 has columns named a, b, and c (and perhaps others) you can make the following query:

```
SELECT a, b + c FROM table1;
```

(assuming that b and c are of a numerical data type). See Section 7.3 for more details.

FROM table1 is a simple kind of table expression: it reads just one table. In general, table expressions can be complex constructs of base tables, joins, and subqueries. But you can also omit the table expression entirely and use the SELECT command as a calculator:

```
SELECT 3 * 4;
```

This is more useful if the expressions in the select list return varying results. For example, you could call a function this way:

```
SELECT random();
```

7.2. Table Expressions

A table expression computes a table. The table expression contains a FROM clause that is optionally followed by WHERE, GROUP BY, and HAVING clauses. Trivial table expressions simply refer to a table on disk, a so-called base table, but more complex expressions can be used to modify or combine base tables in various ways.

The optional WHERE, GROUP BY, and HAVING clauses in the table expression specify a pipeline of successive transformations performed on the table derived in the FROM clause. All these transforma-
tions produce a virtual table that provides the rows that are passed to the select list to compute the output rows of the query.

7.2.1. The FROM Clause

The FROM Clause derives a table from one or more other tables given in a comma-separated table reference list.

FROM table_reference [, table_reference [, ...]]

A table reference can be a table name (possibly schema-qualified), or a derived table such as a sub-query, a JOIN construct, or complex combinations of these. If more than one table reference is listed in the FROM clause, the tables are cross-joined (that is, the Cartesian product of their rows is formed; see below). The result of the FROM list is an intermediate virtual table that can then be subject to transformations by the WHERE, GROUP BY, and HAVING clauses and is finally the result of the overall table expression.

When a table reference names a table that is the parent of a table inheritance hierarchy, the table reference produces rows of not only that table but all of its descendant tables, unless the key word ONLY precedes the table name. However, the reference produces only the columns that appear in the named table — any columns added in subtables are ignored.

Instead of writing ONLY before the table name, you can write * after the table name to explicitly specify that descendant tables are included. There is no real reason to use this syntax any more, because searching descendant tables is now always the default behavior. However, it is supported for compatibility with older releases.

7.2.1.1. Joined Tables

A joined table is a table derived from two other (real or derived) tables according to the rules of the particular join type. Inner, outer, and cross-joins are available. The general syntax of a joined table is

T1 join_type T2 [ join_condition ]

Joins of all types can be chained together, or nested: either or both T1 and T2 can be joined tables. Parentheses can be used around JOIN clauses to control the join order. In the absence of parentheses, JOIN clauses nest left-to-right.

Join Types

Cross join

T1 CROSS JOIN T2

For every possible combination of rows from T1 and T2 (i.e., a Cartesian product), the joined table will contain a row consisting of all columns in T1 followed by all columns in T2. If the tables have N and M rows respectively, the joined table will have N * M rows.

FROM T1 CROSS JOIN T2 is equivalent to FROM T1 INNER JOIN T2 ON TRUE (see below). It is also equivalent to FROM T1, T2.

Note

This latter equivalence does not hold exactly when more than two tables appear, because JOIN binds more tightly than comma. For example FROM T1 CROSS JOIN T2 INNER JOIN T3 ON condition is not the same as FROM
Qualified joins

\[
T1, T2 \text{ INNER JOIN } T3 \text{ ON condition because the condition can reference } T1 \text{ in the first case but not the second.}
\]

The words \text{INNER} and \text{OUTER} are optional in all forms. \text{INNER} is the default; \text{LEFT}, \text{RIGHT}, and \text{FULL} imply an outer join.

The \text{join condition} is specified in the \text{ON} or \text{USING} clause, or implicitly by the word \text{NATURAL}. The join condition determines which rows from the two source tables are considered to “match”, as explained in detail below.

The possible types of qualified join are:

- **\text{INNER JOIN}**
  - For each row \(R1\) of \(T1\), the joined table has a row for each row in \(T2\) that satisfies the join condition with \(R1\).

- **\text{LEFT OUTER JOIN}**
  - First, an inner join is performed. Then, for each row in \(T1\) that does not satisfy the join condition with any row in \(T2\), a joined row is added with null values in columns of \(T2\). Thus, the joined table always has at least one row for each row in \(T1\).

- **\text{RIGHT OUTER JOIN}**
  - First, an inner join is performed. Then, for each row in \(T2\) that does not satisfy the join condition with any row in \(T1\), a joined row is added with null values in columns of \(T1\). This is the converse of a left join: the result table will always have a row for each row in \(T2\).

- **\text{FULL OUTER JOIN}**
  - First, an inner join is performed. Then, for each row in \(T1\) that does not satisfy the join condition with any row in \(T2\), a joined row is added with null values in columns of \(T2\). Also, for each row of \(T2\) that does not satisfy the join condition with any row in \(T1\), a joined row with null values in the columns of \(T1\) is added.

The \text{ON} clause is the most general kind of join condition: it takes a Boolean value expression of the same kind as is used in a \text{WHERE} clause. A pair of rows from \(T1\) and \(T2\) match if the \text{ON} expression evaluates to true.

The \text{USING} clause is a shorthand that allows you to take advantage of the specific situation where both sides of the join use the same name for the joining column(s). It takes a comma-separated list of the shared column names and forms a join condition that includes an equality comparison for each one. For example, joining \(T1\) and \(T2\) with \text{USING} \((a, b)\) produces the join condition \text{ON } T1.a = T2.a AND T1.b = T2.b.

Furthermore, the output of \text{JOIN USING} suppresses redundant columns: there is no need to print both of the matched columns, since they must have equal values. While \text{JOIN ON} produces all columns from \(T1\) followed by all columns from \(T2\), \text{JOIN USING} produces one output column for each of the listed column pairs (in the listed order), followed by any remaining columns from \(T1\), followed by any remaining columns from \(T2\).
Finally, NATURAL is a shorthand form of USING: it forms a USING list consisting of all column names that appear in both input tables. As with USING, these columns appear only once in the output table. If there are no common column names, NATURAL JOIN behaves like JOIN ... ON TRUE, producing a cross-product join.

Note

USING is reasonably safe from column changes in the joined relations since only the listed columns are combined. NATURAL is considerably more risky since any schema changes to either relation that cause a new matching column name to be present will cause the join to combine that new column as well.

To put this together, assume we have tables t1:

```
  num | name
  ----+------
   1  |  a   
   2  |  b   
   3  |  c   
```

and t2:

```
  num | value
  ----+-------
   1  |  xxx  
   3  |  yyy  
   5  |  zzz  
```

then we get the following results for the various joins:

```sql
=> SELECT * FROM t1 CROSS JOIN t2;
  num | name | num | value
  ----+------+-----+-------
   1  |  a   |   1 |  xxx  
   1  |  a   |   3 |  yyy  
   1  |  a   |   5 |  zzz  
   2  |  b   |   1 |  xxx  
   2  |  b   |   3 |  yyy  
   2  |  b   |   5 |  zzz  
   3  |  c   |   1 |  xxx  
   3  |  c   |   3 |  yyy  
   3  |  c   |   5 |  zzz  
(9 rows)
```

```sql
=> SELECT * FROM t1 INNER JOIN t2 ON t1.num = t2.num;
  num | name | num | value
  ----+------+-----+-------
   1  |  a   |   1 |  xxx  
   3  |  c   |   3 |  yyy  
(2 rows)
```

```sql
=> SELECT * FROM t1 INNER JOIN t2 USING (num);
  num | name | value
  ----+------+-------
   1  |  a   |  xxx  
```
The join condition specified with `ON` can also contain conditions that do not relate directly to the join. This can prove useful for some queries but needs to be thought out carefully. For example:

```sql
=> SELECT * FROM t1 LEFT JOIN t2 ON t1.num = t2.num AND t2.value = 'xxx';
num | name | num | value
-----+------+-----+-------
 1   | a    |   1 | xxx
 3   | c    |   3 | yyy
(3 rows)
```

Notice that placing the restriction in the `WHERE` clause produces a different result:
=> SELECT * FROM t1 LEFT JOIN t2 ON t1.num = t2.num WHERE t2.value = 'xxx';

<table>
<thead>
<tr>
<th>num</th>
<th>name</th>
<th>num</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>a</td>
<td>1</td>
<td>xxx</td>
</tr>
</tbody>
</table>

(1 row)

This is because a restriction placed in the ON clause is processed before the join, while a restriction placed in the WHERE clause is processed after the join. That does not matter with inner joins, but it matters a lot with outer joins.

### 7.2.1.2. Table and Column Aliases

A temporary name can be given to tables and complex table references to be used for references to the derived table in the rest of the query. This is called a table alias.

To create a table alias, write

```
FROM table_reference AS alias
```

or

```
FROM table_reference alias
```

The AS key word is optional noise. alias can be any identifier.

A typical application of table aliases is to assign short identifiers to long table names to keep the join clauses readable. For example:

```sql
SELECT * FROM some_very_long_table_name s JOIN another_fairly_long_name a ON s.id = a.num;
```

The alias becomes the new name of the table reference so far as the current query is concerned — it is not allowed to refer to the table by the original name elsewhere in the query. Thus, this is not valid:

```sql
SELECT * FROM my_table AS m WHERE my_table.a > 5;    -- wrong
```

Table aliases are mainly for notational convenience, but it is necessary to use them when joining a table to itself, e.g.:

```sql
SELECT * FROM people AS mother JOIN people AS child ON mother.id = child.mother_id;
```

Additionally, an alias is required if the table reference is a subquery (see Section 7.2.1.3).

Parentheses are used to resolve ambiguities. In the following example, the first statement assigns the alias b to the second instance of my_table, but the second statement assigns the alias to the result of the join:

```sql
SELECT * FROM my_table AS a CROSS JOIN my_table AS b ...
SELECT * FROM (my_table AS a CROSS JOIN my_table) AS b ...
```

Another form of table aliasing gives temporary names to the columns of the table, as well as the table itself:
FROM table_reference [AS] alias ( column1 [, column2 [, ...]])

If fewer column aliases are specified than the actual table has columns, the remaining columns are not renamed. This syntax is especially useful for self-joins or subqueries.

When an alias is applied to the output of a JOIN clause, the alias hides the original name(s) within the JOIN. For example:

SELECT a.* FROM my_table AS a JOIN your_table AS b ON ...

is valid SQL, but:

SELECT a.* FROM (my_table AS a JOIN your_table AS b ON ...) AS c

is not valid; the table alias a is not visible outside the alias c.

7.2.1.3. Subqueries

Subqueries specifying a derived table must be enclosed in parentheses and must be assigned a table alias name (as in Section 7.2.1.2). For example:

FROM (SELECT * FROM table1) AS alias_name

This example is equivalent to FROM table1 AS alias_name. More interesting cases, which cannot be reduced to a plain join, arise when the subquery involves grouping or aggregation.

A subquery can also be a VALUES list:

FROM (VALUES ('anne', 'smith'), ('bob', 'jones'), ('joe', 'blow')) AS names(first, last)

Again, a table alias is required. Assigning alias names to the columns of the VALUES list is optional, but is good practice. For more information see Section 7.7.

7.2.1.4. Table Functions

Table functions are functions that produce a set of rows, made up of either base data types (scalar types) or composite data types (table rows). They are used like a table, view, or subquery in the FROM clause of a query. Columns returned by table functions can be included in SELECT, JOIN, or WHERE clauses in the same manner as columns of a table, view, or subquery.

Table functions may also be combined using the ROWS FROM syntax, with the results returned in parallel columns; the number of result rows in this case is that of the largest function result, with smaller results padded with null values to match.

function_call [WITH ORDINALITY] [[AS] table_alias [(column_alias [, ...])]]
ROWS FROM( function_call [, ... ] ) [WITH ORDINALITY] [[AS] table_alias [(column_alias [, ...])]]

If the WITH ORDINALITY clause is specified, an additional column of type bigint will be added to the function result columns. This column numbers the rows of the function result set, starting from 1. (This is a generalization of the SQL-standard syntax for UNNEST ... WITH ORDINALITY.)
By default, the ordinal column is called `ordinality`, but a different column name can be assigned to it using an `AS` clause.

The special table function `UNNEST` may be called with any number of array parameters, and it returns a corresponding number of columns, as if `UNNEST` (Section 9.18) had been called on each parameter separately and combined using the `ROWS FROM()` construct.

```
UNNEST( array_expression [, ... ] ) [WITH ORDINALITY]
  [[AS] table_alias [(column_alias [, ... ])]]
```

If no `table_alias` is specified, the function name is used as the table name; in the case of a `ROWS FROM()` construct, the first function's name is used.

If column aliases are not supplied, then for a function returning a base data type, the column name is also the same as the function name. For a function returning a composite type, the result columns get the names of the individual attributes of the type.

Some examples:

```
CREATE TABLE foo (fooid int, foosubid int, fooname text);

CREATE FUNCTION getfoo(int) RETURNS SETOF foo AS $$
  SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;

SELECT * FROM getfoo(1) AS t1;

SELECT * FROM foo
  WHERE foosubid IN
    (SELECT foosubid
     FROM getfoo(foo.fooid) z
     WHERE z.fooid = foo.fooid);

CREATE VIEW vw_getfoo AS SELECT * FROM getfoo(1);

SELECT * FROM vw_getfoo;
```

In some cases it is useful to define table functions that can return different column sets depending on how they are invoked. To support this, the table function can be declared as returning the pseudo-type `record`. When such a function is used in a query, the expected row structure must be specified in the query itself, so that the system can know how to parse and plan the query. This syntax looks like:

```
function_call [AS] alias (column_definition [, ... ])
function_call AS [alias] (column_definition [, ... ])
ROWS FROM( ... function_call AS (column_definition [, ... ])
  [, ... ] )
```

When not using the `ROWS FROM()` syntax, the `column_definition` list replaces the column alias list that could otherwise be attached to the `FROM` item; the names in the column definitions serve as column aliases. When using the `ROWS FROM()` syntax, a `column_definition` list can be attached to each member function separately; or if there is only one member function and no `WITH ORDINALITY` clause, a `column_definition` list can be written in place of a column alias list following `ROWS FROM()`.

Consider this example:
Queries

```sql
SELECT *
FROM dblink('dbname=mydb', 'SELECT proname, prosrc FROM pg_proc')
AS t1(proname name, prosrc text)
WHERE proname LIKE 'bytea%';
```

The `dblink` function (part of the `dblink` module) executes a remote query. It is declared to return record since it might be used for any kind of query. The actual column set must be specified in the calling query so that the parser knows, for example, what * should expand to.

### 7.2.1.5. LATERAL Subqueries

Subqueries appearing in `FROM` can be preceded by the key word `LATERAL`. This allows them to reference columns provided by preceding `FROM` items. (Without `LATERAL`, each subquery is evaluated independently and so cannot cross-reference any other `FROM` item.)

Table functions appearing in `FROM` can also be preceded by the key word `LATERAL`, but for functions the key word is optional; the function's arguments can contain references to columns provided by preceding `FROM` items in any case.

A `LATERAL` item can appear at top level in the `FROM` list, or within a `JOIN` tree. In the latter case it can also refer to any items that are on the left-hand side of a `JOIN` that it is on the right-hand side of.

When a `FROM` item contains `LATERAL` cross-references, evaluation proceeds as follows: for each row of the `FROM` item providing the cross-referenced column(s), or set of rows of multiple `FROM` items providing the columns, the `LATERAL` item is evaluated using that row or row set's values of the columns. The resulting row(s) are joined as usual with the rows they were computed from. This is repeated for each row or set of rows from the column source table(s).

A trivial example of `LATERAL` is

```sql
SELECT * FROM foo, LATERAL (SELECT * FROM bar WHERE bar.id = foo.bar_id) ss;
```

This is not especially useful since it has exactly the same result as the more conventional

```sql
SELECT * FROM foo, bar WHERE bar.id = foo.bar_id;
```

`LATERAL` is primarily useful when the cross-referenced column is necessary for computing the row(s) to be joined. A common application is providing an argument value for a set-returning function. For example, supposing that `vertices(polygon)` returns the set of vertices of a polygon, we could identify close-together vertices of polygons stored in a table with:

```sql
SELECT p1.id, p2.id, v1, v2
FROM polygons p1, polygons p2,
    LATERAL vertices(p1.poly) v1,
    LATERAL vertices(p2.poly) v2
WHERE (v1 <-> v2) < 10 AND p1.id != p2.id;
```

This query could also be written

```sql
SELECT p1.id, p2.id, v1, v2
FROM polygons p1 CROSS JOIN LATERAL vertices(p1.poly) v1,
    polygons p2 CROSS JOIN LATERAL vertices(p2.poly) v2
WHERE (v1 <-> v2) < 10 AND p1.id != p2.id;
```
or in several other equivalent formulations. (As already mentioned, the LATERAL key word is unnecessary in this example, but we use it for clarity.)

It is often particularly handy to LEFT JOIN to a LATERAL subquery, so that source rows will appear in the result even if the LATERAL subquery produces no rows for them. For example, if `get_product_names()` returns the names of products made by a manufacturer, but some manufacturers in our table currently produce no products, we could find out which ones those are like this:

```sql
SELECT m.name
FROM manufacturers m LEFT JOIN LATERAL get_product_names(m.id) pname ON true
WHERE pname IS NULL;
```

### 7.2.2. The **WHERE** Clause

The syntax of the **WHERE** Clause is:

```
WHERE search_condition
```

where `search_condition` is any value expression (see Section 4.2) that returns a value of type boolean.

After the processing of the FROM clause is done, each row of the derived virtual table is checked against the search condition. If the result of the condition is true, the row is kept in the output table, otherwise (i.e., if the result is false or null) it is discarded. The search condition typically references at least one column of the table generated in the FROM clause; this is not required, but otherwise the WHERE clause will be fairly useless.

**Note**

The join condition of an inner join can be written either in the WHERE clause or in the JOIN clause. For example, these table expressions are equivalent:

```sql
FROM a, b WHERE a.id = b.id AND b.val > 5
```

and:

```sql
FROM a INNER JOIN b ON (a.id = b.id) WHERE b.val > 5
```

or perhaps even:

```sql
FROM a NATURAL JOIN b WHERE b.val > 5
```

Which one of these you use is mainly a matter of style. The JOIN syntax in the FROM clause is probably not as portable to other SQL database management systems, even though it is in the SQL standard. For outer joins there is no choice: they must be done in the FROM clause. The ON or USING clause of an outer join is not equivalent to a WHERE condition, because it results in the addition of rows (for unmatched input rows) as well as the removal of rows in the final result.

Here are some examples of **WHERE** clauses:
### Queries

```sql
SELECT ... FROM fdt WHERE c1 > 5
SELECT ... FROM fdt WHERE c1 IN (1, 2, 3)
SELECT ... FROM fdt WHERE c1 IN (SELECT c1 FROM t2)
SELECT ... FROM fdt WHERE c1 IN (SELECT c3 FROM t2 WHERE c2 = fdt.c1 + 10)
SELECT ... FROM fdt WHERE c1 BETWEEN (SELECT c3 FROM t2 WHERE c2 = fdt.c1 + 10) AND 100
SELECT ... FROM fdt WHERE EXISTS (SELECT c1 FROM t2 WHERE c2 > fdt.c1)
```

`fdt` is the table derived in the FROM clause. Rows that do not meet the search condition of the WHERE clause are eliminated from `fdt`. Notice the use of scalar subqueries as value expressions. Just like any other query, the subqueries can employ complex table expressions. Notice also how `fdt` is referenced in the subqueries. Qualifying `c1` as `fdt.c1` is only necessary if `c1` is also the name of a column in the derived input table of the subquery. But qualifying the column name adds clarity even when it is not needed. This example shows how the column naming scope of an outer query extends into its inner queries.

### 7.2.3. The `GROUP BY` and `HAVING` Clauses

After passing the WHERE filter, the derived input table might be subject to grouping, using the GROUP BY clause, and elimination of group rows using the HAVING clause.

```sql
SELECT select_list
FROM ...
[WHERE ...]
GROUP BY grouping_column_reference
[, grouping_column_reference]...
```

The `GROUP BY` Clause is used to group together those rows in a table that have the same values in all the columns listed. The order in which the columns are listed does not matter. The effect is to combine each set of rows having common values into one group row that represents all rows in the group. This is done to eliminate redundancy in the output and/or compute aggregates that apply to these groups. For instance:

```sql
=> SELECT * FROM test1;
  x | y
----+---
  a | 3
  c | 2
  b | 5
  a | 1
(4 rows)

=> SELECT x FROM test1 GROUP BY x;
  x
---
  a
  b
  c
(3 rows)
```
In the second query, we could not have written `SELECT * FROM test1 GROUP BY x`, because there is no single value for the column `y` that could be associated with each group. The grouped-by columns can be referenced in the select list since they have a single value in each group.

In general, if a table is grouped, columns that are not listed in `GROUP BY` cannot be referenced except in aggregate expressions. An example with aggregate expressions is:

```sql
=> SELECT x, sum(y) FROM test1 GROUP BY x;
```

<table>
<thead>
<tr>
<th>x</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>4</td>
</tr>
<tr>
<td>b</td>
<td>5</td>
</tr>
<tr>
<td>c</td>
<td>2</td>
</tr>
</tbody>
</table>

(3 rows)

Here `sum` is an aggregate function that computes a single value over the entire group. More information about the available aggregate functions can be found in Section 9.20.

**Tip**

Grouping without aggregate expressions effectively calculates the set of distinct values in a column. This can also be achieved using the `DISTINCT` clause (see Section 7.3.3).

Here is another example: it calculates the total sales for each product (rather than the total sales of all products):

```sql
SELECT product_id, p.name, (sum(s.units) * p.price) AS sales
FROM products p LEFT JOIN sales s USING (product_id)
GROUP BY product_id, p.name, p.price;
```

In this example, the columns `product_id`, `p.name`, and `p.price` must be in the `GROUP BY` clause since they are referenced in the query select list (but see below). The column `s.units` does not have to be in the `GROUP BY` list since it is only used in an aggregate expression (`sum(...)`) which represents the sales of a product. For each product, the query returns a summary row about all sales of the product.

If the products table is set up so that, say, `product_id` is the primary key, then it would be enough to group by `product_id` in the above example, since name and price would be *functionally dependent* on the product ID, and so there would be no ambiguity about which name and price value to return for each product ID group.

In strict SQL, `GROUP BY` can only group by columns of the source table but PostgreSQL extends this to also allow `GROUP BY` to group by columns in the select list. Grouping by value expressions instead of simple column names is also allowed.

If a table has been grouped using `GROUP BY`, but only certain groups are of interest, the `HAVING` clause can be used, much like a `WHERE` clause, to eliminate groups from the result. The syntax is:

```sql
SELECT select_list FROM ... [WHERE ...] GROUP BY ...
HAVING boolean_expression
```

Expressions in the `HAVING` clause can refer both to grouped expressions and to ungrouped expressions (which necessarily involve an aggregate function).
Example:

```sql
=> SELECT x, sum(y) FROM test1 GROUP BY x HAVING sum(y) > 3;
```
```
<table>
<thead>
<tr>
<th>x</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>4</td>
</tr>
<tr>
<td>b</td>
<td>5</td>
</tr>
</tbody>
</table>
```

(2 rows)

```sql
=> SELECT x, sum(y) FROM test1 GROUP BY x HAVING x < 'c';
```
```
<table>
<thead>
<tr>
<th>x</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>4</td>
</tr>
<tr>
<td>b</td>
<td>5</td>
</tr>
</tbody>
</table>
```

(2 rows)

Again, a more realistic example:

```sql
SELECT product_id, p.name, (sum(s.units) * (p.price - p.cost)) AS profit
FROM products p LEFT JOIN sales s USING (product_id)
WHERE s.date > CURRENT_DATE - INTERVAL '4 weeks'
GROUP BY product_id, p.name, p.price, p.cost
HAVING sum(p.price * s.units) > 5000;
```

In the example above, the WHERE clause is selecting rows by a column that is not grouped (the expression is only true for sales during the last four weeks), while the HAVING clause restricts the output to groups with total gross sales over 5000. Note that the aggregate expressions do not necessarily need to be the same in all parts of the query.

If a query contains aggregate function calls, but no GROUP BY clause, grouping still occurs: the result is a single group row (or perhaps no rows at all, if the single row is then eliminated by HAVING). The same is true if it contains a HAVING clause, even without any aggregate function calls or GROUP BY clause.

### 7.2.4. GROUPING SETS, CUBE, and ROLLUP

More complex grouping operations than those described above are possible using the concept of grouping sets. The data selected by the FROM and WHERE clauses is grouped separately by each specified grouping set, aggregates computed for each group just as for simple GROUP BY clauses, and then the results returned. For example:

```sql
=> SELECT * FROM items_sold;
```
```
<table>
<thead>
<tr>
<th>brand</th>
<th>size</th>
<th>sales</th>
</tr>
</thead>
<tbody>
<tr>
<td>Foo</td>
<td>L</td>
<td>10</td>
</tr>
<tr>
<td>Foo</td>
<td>M</td>
<td>20</td>
</tr>
<tr>
<td>Bar</td>
<td>M</td>
<td>15</td>
</tr>
<tr>
<td>Bar</td>
<td>L</td>
<td>5</td>
</tr>
</tbody>
</table>
```

(4 rows)

```sql
=> SELECT brand, size, sum(sales) FROM items_sold GROUP BY GROUPING SETS ((brand), (size), ());
```
```
<table>
<thead>
<tr>
<th>brand</th>
<th>size</th>
<th>sum</th>
</tr>
</thead>
</table>
```

(2 rows)
Each sublist of `GROUPING SETS` may specify zero or more columns or expressions and is interpreted the same way as though it were directly in the `GROUP BY` clause. An empty grouping set means that all rows are aggregated down to a single group (which is output even if no input rows were present), as described above for the case of aggregate functions with no `GROUP BY` clause.

References to the grouping columns or expressions are replaced by null values in result rows for grouping sets in which those columns do not appear. To distinguish which grouping a particular output row resulted from, see Table 9.56.

A shorthand notation is provided for specifying two common types of grouping set. A clause of the form

```
ROLLUP ( e1, e2, e3, ... )
```

represents the given list of expressions and all prefixes of the list including the empty list; thus it is equivalent to

```
GROUPING SETS (  
  ( e1, e2, e3, ... ),  
  ...,  
  ( e1, e2 ),  
  ( e1 ),  
  ( ),
)
```

This is commonly used for analysis over hierarchical data; e.g. total salary by department, division, and company-wide total.

A clause of the form

```
CUBE ( e1, e2, ... )
```

represents the given list and all of its possible subsets (i.e. the power set). Thus

```
CUBE ( a, b, c )
```

is equivalent to

```
GROUPING SETS (  
  ( a, b, c ),  
  ( a, b ),  
  ( a, c ),  
  ( a ),  
  ( b, c ),  
  ( b ),  
  ( c ),  
  ( )
)
```
The individual elements of a `CUBE` or `ROLLUP` clause may be either individual expressions, or sublists of elements in parentheses. In the latter case, the sublists are treated as single units for the purposes of generating the individual grouping sets. For example:

```
CUBE ( (a, b), (c, d) )
```

is equivalent to

```
GROUPING SETS ( ( a, b, c, d ), ( a, b ), ( c, d ), ( ) )
```

and

```
ROLLUP ( a, (b, c), d )
```

is equivalent to

```
GROUPING SETS ( ( a, b, c, d ), ( a, b, c ), ( a ), ( ) )
```

The `CUBE` and `ROLLUP` constructs can be used either directly in the `GROUP BY` clause, or nested inside a `GROUPING SETS` clause. If one `GROUPING SETS` clause is nested inside another, the effect is the same as if all the elements of the inner clause had been written directly in the outer clause.

If multiple grouping items are specified in a single `GROUP BY` clause, then the final list of grouping sets is the cross product of the individual items. For example:

```
GROUP BY a, CUBE (b, c), GROUPING SETS ((d), (e))
```

is equivalent to

```
GROUP BY GROUPING SETS ( (a, b, c, d), (a, b, c, e), (a, b, d), (a, b, e), (a, c, d), (a, c, e), (a, d), (a, e) )
```

Note

The construct `(a, b)` is normally recognized in expressions as a row constructor. Within the `GROUP BY` clause, this does not apply at the top levels of expressions, and `(a, b)` is parsed as a list of expressions as described above. If for some reason you need a row constructor in a grouping expression, use `ROW(a, b)`.
7.2.5. Window Function Processing

If the query contains any window functions (see Section 3.5, Section 9.21 and Section 4.2.8), these functions are evaluated after any grouping, aggregation, and HAVING filtering is performed. That is, if the query uses any aggregates, GROUP BY, or HAVING, then the rows seen by the window functions are the group rows instead of the original table rows from FROM/WHERE.

When multiple window functions are used, all the window functions having syntactically equivalent PARTITION BY and ORDER BY clauses in their window definitions are guaranteed to be evaluated in a single pass over the data. Therefore they will see the same sort ordering, even if the ORDER BY does not uniquely determine an ordering. However, no guarantees are made about the evaluation of functions having different PARTITION BY or ORDER BY specifications. (In such cases a sort step is typically required between the passes of window function evaluations, and the sort is not guaranteed to preserve ordering of rows that its ORDER BY sees as equivalent.)

Currently, window functions always require presorted data, and so the query output will be ordered according to one or another of the window functions' PARTITION BY/ORDER BY clauses. It is not recommended to rely on this, however. Use an explicit top-level ORDER BY clause if you want to be sure the results are sorted in a particular way.

7.3. Select Lists

As shown in the previous section, the table expression in the SELECT command constructs an intermediate virtual table by possibly combining tables, views, eliminating rows, grouping, etc. This table is finally passed on to processing by the select list. The select list determines which columns of the intermediate table are actually output.

7.3.1. Select-List Items

The simplest kind of select list is * which emits all columns that the table expression produces. Otherwise, a select list is a comma-separated list of value expressions (as defined in Section 4.2). For instance, it could be a list of column names:

```
SELECT a, b, c FROM ...
```

The columns names a, b, and c are either the actual names of the columns of tables referenced in the FROM clause, or the aliases given to them as explained in Section 7.2.1.2. The name space available in the select list is the same as in the WHERE clause, unless grouping is used, in which case it is the same as in the HAVING clause.

If more than one table has a column of the same name, the table name must also be given, as in:

```
SELECT tbl1.a, tbl2.a, tbl1.b FROM ...
```

When working with multiple tables, it can also be useful to ask for all the columns of a particular table:

```
SELECT tbl1.*, tbl2.a FROM ...
```

See Section 8.16.5 for more about the table_name.* notation.

If an arbitrary value expression is used in the select list, it conceptually adds a new virtual column to the returned table. The value expression is evaluated once for each result row, with the row's values substituted for any column references. But the expressions in the select list do not have to reference any columns in the table expression of the FROM clause; they can be constant arithmetic expressions, for instance.
7.3.2. Column Labels

The entries in the select list can be assigned names for subsequent processing, such as for use in an ORDER BY clause or for display by the client application. For example:

```sql
SELECT a AS value, b + c AS sum FROM ...
```

If no output column name is specified using AS, the system assigns a default column name. For simple column references, this is the name of the referenced column. For function calls, this is the name of the function. For complex expressions, the system will generate a generic name.

The AS keyword is optional, but only if the new column name does not match any PostgreSQL keyword (see Appendix C). To avoid an accidental match to a keyword, you can double-quote the column name. For example, VALUE is a keyword, so this does not work:

```sql
SELECT a value, b + c AS sum FROM ...
```

but this does:

```sql
SELECT a "value", b + c AS sum FROM ...
```

For protection against possible future keyword additions, it is recommended that you always either write AS or double-quote the output column name.

**Note**

The naming of output columns here is different from that done in the FROM clause (see Section 7.2.1.2). It is possible to rename the same column twice, but the name assigned in the select list is the one that will be passed on.

7.3.3. DISTINCT

After the select list has been processed, the result table can optionally be subject to the elimination of duplicate rows. The DISTINCT key word is written directly after SELECT to specify this:

```sql
SELECT DISTINCT select_list ...
```

(Instead of DISTINCT the key word ALL can be used to specify the default behavior of retaining all rows.)

Obviously, two rows are considered distinct if they differ in at least one column value. Null values are considered equal in this comparison.

Alternatively, an arbitrary expression can determine what rows are to be considered distinct:

```sql
SELECT DISTINCT ON (expression [, expression ...]) select_list ...
```

Here expression is an arbitrary value expression that is evaluated for all rows. A set of rows for which all the expressions are equal are considered duplicates, and only the first row of the set is kept in the output. Note that the “first row” of a set is unpredictable unless the query is sorted on enough columns to guarantee a unique ordering of the rows arriving at the DISTINCT filter. (DISTINCT ON processing occurs after ORDER BY sorting.)
The `DISTINCT ON` clause is not part of the SQL standard and is sometimes considered bad style because of the potentially indeterminate nature of its results. With judicious use of `GROUP BY` and subqueries in `FROM`, this construct can be avoided, but it is often the most convenient alternative.

### 7.4. Combining Queries

The results of two queries can be combined using the set operations union, intersection, and difference. The syntax is

```sql
query1 UNION [ALL] query2
query1 INTERSECT [ALL] query2
query1 EXCEPT [ALL] query2
```

`query1` and `query2` are queries that can use any of the features discussed up to this point. Set operations can also be nested and chained, for example

```sql
query1 UNION query2 UNION query3
```

which is executed as:

```sql
(query1 UNION query2) UNION query3
```

`UNION` effectively appends the result of `query2` to the result of `query1` (although there is no guarantee that this is the order in which the rows are actually returned). Furthermore, it eliminates duplicate rows from its result, in the same way as `DISTINCT`, unless `UNION ALL` is used.

`INTERSECT` returns all rows that are both in the result of `query1` and in the result of `query2`. Duplicate rows are eliminated unless `INTERSECT ALL` is used.

`EXCEPT` returns all rows that are in the result of `query1` but not in the result of `query2`. (This is sometimes called the difference between two queries.) Again, duplicates are eliminated unless `EXCEPT ALL` is used.

In order to calculate the union, intersection, or difference of two queries, the two queries must be “union compatible”, which means that they return the same number of columns and the corresponding columns have compatible data types, as described in Section 10.5.

### 7.5. Sorting Rows

After a query has produced an output table (after the select list has been processed) it can optionally be sorted. If sorting is not chosen, the rows will be returned in an unspecified order. The actual order in that case will depend on the scan and join plan types and the order on disk, but it must not be relied on. A particular output ordering can only be guaranteed if the sort step is explicitly chosen.

The `ORDER BY` clause specifies the sort order:

```sql
SELECT select_list
    FROM table_expression
    ORDER BY sort_expression1 [ASC | DESC] [NULLS { FIRST | LAST }]
            [, sort_expression2 [ASC | DESC] [NULLS { FIRST | LAST }]] ...
```

The sort expression(s) can be any expression that would be valid in the query’s select list. An example is:
SELECT a, b FROM table1 ORDER BY a + b, c;

When more than one expression is specified, the later values are used to sort rows that are equal according to the earlier values. Each expression can be followed by an optional **ASC** or **DESC** keyword to set the sort direction to ascending or descending. **ASC** order is the default. Ascending order puts smaller values first, where “smaller” is defined in terms of the `<` operator. Similarly, descending order is determined with the `>` operator.¹

The **NULLS FIRST** and **NULLS LAST** options can be used to determine whether nulls appear before or after non-null values in the sort ordering. By default, null values sort as if larger than any non-null value; that is, **NULLS FIRST** is the default for **DESC** order, and **NULLS LAST** otherwise.

Note that the ordering options are considered independently for each sort column. For example `ORDER BY x, y DESC` means `ORDER BY x ASC, y DESC`, which is not the same as `ORDER BY x DESC, y DESC`.

A **sort_expression** can also be the column label or number of an output column, as in:

- `SELECT a + b AS sum, c FROM table1 ORDER BY sum;`
- `SELECT a, max(b) FROM table1 GROUP BY a ORDER BY 1;`

both of which sort by the first output column. Note that an output column name has to stand alone, that is, it cannot be used in an expression — for example, this is **not** correct:

- `SELECT a + b AS sum, c FROM table1 ORDER BY sum + c;` -- wrong

This restriction is made to reduce ambiguity. There is still ambiguity if an **ORDER BY** item is a simple name that could match either an output column name or a column from the table expression. The output column is used in such cases. This would only cause confusion if you use **AS** to rename an output column to match some other table column’s name.

**ORDER BY** can be applied to the result of a **UNION**, **INTERSECT**, or **EXCEPT** combination, but in this case it is only permitted to sort by output column names or numbers, not by expressions.

### 7.6. LIMIT and OFFSET

**LIMIT** and **OFFSET** allow you to retrieve just a portion of the rows that are generated by the rest of the query:

```sql
SELECT select_list
FROM table_expression
[ ORDER BY ... ]
[ LIMIT { number | ALL } ] [ OFFSET number ]
```

If a limit count is given, no more than that many rows will be returned (but possibly fewer, if the query itself yields fewer rows). **LIMIT ALL** is the same as omitting the **LIMIT** clause, as is **LIMIT** with a NULL argument.

**OFFSET** says to skip that many rows before beginning to return rows. **OFFSET 0** is the same as omitting the **OFFSET** clause, as is **OFFSET** with a NULL argument.

If both **OFFSET** and **LIMIT** appear, then **OFFSET** rows are skipped before starting to count the **LIMIT** rows that are returned.

---

¹ Actually, PostgreSQL uses the **default B-tree operator class** for the expression’s data type to determine the sort ordering for **ASC** and **DESC**. Conventionally, data types will be set up so that the `<` and `>` operators correspond to this sort ordering, but a user-defined data type’s designer could choose to do something different.
When using LIMIT, it is important to use an ORDER BY clause that constrains the result rows into a unique order. Otherwise you will get an unpredictable subset of the query's rows. You might be asking for the tenth through twentieth rows, but tenth through twentieth in what ordering? The ordering is unknown, unless you specified ORDER BY.

The query optimizer takes LIMIT into account when generating query plans, so you are very likely to get different plans (yielding different row orders) depending on what you give for LIMIT and OFFSET. Thus, using different LIMIT/OFFSET values to select different subsets of a query result will give inconsistent results unless you enforce a predictable result ordering with ORDER BY. This is not a bug; it is an inherent consequence of the fact that SQL does not promise to deliver the results of a query in any particular order unless ORDER BY is used to constrain the order.

The rows skipped by an OFFSET clause still have to be computed inside the server; therefore a large OFFSET might be inefficient.

7.7. VALUES Lists

VALUES provides a way to generate a “constant table” that can be used in a query without having to actually create and populate a table on-disk. The syntax is

VALUES ( expression [, ...] ) [, ...]

Each parenthesized list of expressions generates a row in the table. The lists must all have the same number of elements (i.e., the number of columns in the table), and corresponding entries in each list must have compatible data types. The actual data type assigned to each column of the result is determined using the same rules as for UNION (see Section 10.5).

As an example:

VALUES (1, 'one'), (2, 'two'), (3, 'three');

will return a table of two columns and three rows. It's effectively equivalent to:

```
SELECT 1 AS column1, 'one' AS column2
UNION ALL
SELECT 2, 'two'
UNION ALL
SELECT 3, 'three';
```

By default, PostgreSQL assigns the names column1, column2, etc. to the columns of a VALUES table. The column names are not specified by the SQL standard and different database systems do it differently, so it's usually better to override the default names with a table alias list, like this:

```
=> SELECT * FROM (VALUES (1, 'one'), (2, 'two'), (3, 'three')) AS t (num, letter);
num | letter
-----+--------
 1 | one
 2 | two
 3 | three
(3 rows)
```

Syntactically, VALUES followed by expression lists is treated as equivalent to:
SELECT select_list FROM table_expression

and can appear anywhere a SELECT can. For example, you can use it as part of a UNION, or attach a sort_specification (ORDER BY, LIMIT, and/or OFFSET) to it. VALUES is most commonly used as the data source in an INSERT command, and next most commonly as a subquery.

For more information see VALUES.

### 7.8. WITH Queries (Common Table Expressions)

WITH provides a way to write auxiliary statements for use in a larger query. These statements, which are often referred to as Common Table Expressions or CTEs, can be thought of as defining temporary tables that exist just for one query. Each auxiliary statement in a WITH clause can be a SELECT, INSERT, UPDATE, or DELETE; and the WITH clause itself is attached to a primary statement that can also be a SELECT, INSERT, UPDATE, or DELETE.

#### 7.8.1. SELECT in WITH

The basic value of SELECT in WITH is to break down complicated queries into simpler parts. An example is:

```sql
WITH regional_sales AS (  
    SELECT region, SUM(amount) AS total_sales  
    FROM orders  
    GROUP BY region  
), top_regions AS (  
    SELECT region  
    FROM regional_sales  
    WHERE total_sales > (SELECT SUM(total_sales)/10 FROM regional_sales)  
)  
SELECT region,  
    product,  
    SUM(quantity) AS product_units,  
    SUM(amount) AS product_sales  
FROM orders  
WHERE region IN (SELECT region FROM top_regions)  
GROUP BY region, product;
```

which displays per-product sales totals in only the top sales regions. The WITH clause defines two auxiliary statements named regional_sales and top_regions, where the output of regional_sales is used in top_regions and the output of top_regions is used in the primary SELECT query. This example could have been written without WITH, but we'd have needed two levels of nested sub-SELECTs. It's a bit easier to follow this way.

The optional RECURSIVE modifier changes WITH from a mere syntactic convenience into a feature that accomplishes things not otherwise possible in standard SQL. Using RECURSIVE, a WITH query can refer to its own output. A very simple example is this query to sum the integers from 1 through 100:

```sql
WITH RECURSIVE t(n) AS (  
    VALUES (1)  
    UNION ALL  
    SELECT n+1 FROM t WHERE n < 100  
)
```
The general form of a recursive WITH query is always a non-recursive term, then UNION (or UNION ALL), then a recursive term, where only the recursive term can contain a reference to the query's own output. Such a query is executed as follows:

**Recursive Query Evaluation**

1. Evaluate the non-recursive term. For UNION (but not UNION ALL), discard duplicate rows. Include all remaining rows in the result of the recursive query, and also place them in a temporary *working table*.

2. So long as the working table is not empty, repeat these steps:
   a. Evaluate the recursive term, substituting the current contents of the working table for the recursive self-reference. For UNION (but not UNION ALL), discard duplicate rows and rows that duplicate any previous result row. Include all remaining rows in the result of the recursive query, and also place them in a temporary *intermediate table*.
   b. Replace the contents of the working table with the contents of the intermediate table, then empty the intermediate table.

**Note**

Strictly speaking, this process is iteration not recursion, but RECURSIVE is the terminology chosen by the SQL standards committee.

In the example above, the working table has just a single row in each step, and it takes on the values from 1 through 100 in successive steps. In the 100th step, there is no output because of the WHERE clause, and so the query terminates.

Recursive queries are typically used to deal with hierarchical or tree-structured data. A useful example is this query to find all the direct and indirect sub-parts of a product, given only a table that shows immediate inclusions:

```sql
WITH RECURSIVE included_parts(sub_part, part, quantity) AS (
    SELECT sub_part, part, quantity FROM parts WHERE part = 'our_product'
    UNION ALL
    SELECT p.sub_part, p.part, p.quantity
    FROM included_parts pr, parts p
    WHERE p.part = pr.sub_part
)
SELECT sub_part, SUM(quantity) as total_quantity
FROM included_parts
GROUP BY sub_part
```

When working with recursive queries it is important to be sure that the recursive part of the query will eventually return no tuples, or else the query will loop indefinitely. Sometimes, using UNION instead of UNION ALL can accomplish this by discarding rows that duplicate previous output rows. However, often a cycle does not involve output rows that are completely duplicate: it may be necessary to check just one or a few fields to see if the same point has been reached before. The standard method for handling such situations is to compute an array of the already-visited values. For example, consider the following query that searches a table *graph* using a *link* field:

```sql
WITH RECURSIVE search_graph(id, link, data, depth) AS ( 
```
queries

SELECT g.id, g.link, g.data, 1
FROM graph g
UNION ALL
SELECT g.id, g.link, g.data, sg.depth + 1
FROM graph g, search_graph sg
WHERE g.id = sg.link
)
SELECT * FROM search_graph;

This query will loop if the link relationships contain cycles. Because we require a “depth” output, just changing UNION ALL to UNION would not eliminate the looping. Instead we need to recognize whether we have reached the same row again while following a particular path of links. We add two columns path and cycle to the loop-prone query:

WITH RECURSIVE search_graph(id, link, data, depth, path, cycle) AS
(
  SELECT g.id, g.link, g.data, 1,
      ARRAY[g.id],
      false
  FROM graph g
  UNION ALL
  SELECT g.id, g.link, g.data, sg.depth + 1,
      path || g.id,
      g.id = ANY(path)
  FROM graph g, search_graph sg
  WHERE g.id = sg.link AND NOT cycle
)
SELECT * FROM search_graph;

Aside from preventing cycles, the array value is often useful in its own right as representing the “path” taken to reach any particular row.

In the general case where more than one field needs to be checked to recognize a cycle, use an array of rows. For example, if we needed to compare fields f1 and f2:

WITH RECURSIVE search_graph(id, link, data, depth, path, cycle) AS
(
  SELECT g.id, g.link, g.data, 1,
      ARRAY[ROW(g.f1, g.f2)],
      false
  FROM graph g
  UNION ALL
  SELECT g.id, g.link, g.data, sg.depth + 1,
      path || ROW(g.f1, g.f2),
      ROW(g.f1, g.f2) = ANY(path)
  FROM graph g, search_graph sg
  WHERE g.id = sg.link AND NOT cycle
)
SELECT * FROM search_graph;

Tip

Omit the ROW() syntax in the common case where only one field needs to be checked to recognize a cycle. This allows a simple array rather than a composite-type array to be used, gaining efficiency.
The recursive query evaluation algorithm produces its output in breadth-first search order. You can display the results in depth-first search order by making the outer query ORDER BY a “path” column constructed in this way.

A helpful trick for testing queries when you are not certain if they might loop is to place a LIMIT in the parent query. For example, this query would loop forever without the LIMIT:

```sql
WITH RECURSIVE t(n) AS (
    SELECT 1
    UNION ALL
    SELECT n+1 FROM t
)
SELECT n FROM t LIMIT 100;
```

This works because PostgreSQL’s implementation evaluates only as many rows of a WITH query as are actually fetched by the parent query. Using this trick in production is not recommended, because other systems might work differently. Also, it usually won’t work if you make the outer query sort the recursive query’s results or join them to some other table, because in such cases the outer query will usually try to fetch all of the WITH query’s output anyway.

A useful property of WITH queries is that they are evaluated only once per execution of the parent query, even if they are referred to more than once by the parent query or sibling WITH queries. Thus, expensive calculations that are needed in multiple places can be placed within a WITH query to avoid redundant work. Another possible application is to prevent unwanted multiple evaluations of functions with side-effects. However, the other side of this coin is that the optimizer is less able to push restrictions from the parent query down into a WITH query than an ordinary subquery. The WITH query will generally be evaluated as written, without suppression of rows that the parent query might discard afterwards. (But, as mentioned above, evaluation might stop early if the reference(s) to the query demand only a limited number of rows.)

The examples above only show WITH being used with SELECT, but it can be attached in the same way to INSERT, UPDATE, or DELETE. In each case it effectively provides temporary table(s) that can be referred to in the main command.

### 7.8.2. Data-Modifying Statements in WITH

You can use data-modifying statements (INSERT, UPDATE, or DELETE) in WITH. This allows you to perform several different operations in the same query. An example is:

```sql
WITH moved_rows AS (
    DELETE FROM products
    WHERE
        "date" >= '2010-10-01' AND
        "date" < '2010-11-01'
    RETURNING *
)
INSERT INTO products_log
SELECT * FROM moved_rows;
```

This query effectively moves rows from `products` to `products_log`. The DELETE in WITH deletes the specified rows from `products`, returning their contents by means of its RETURNING clause; and then the primary query reads that output and inserts it into `products_log`.
A fine point of the above example is that the WITH clause is attached to the INSERT, not the sub-
SELECT within the INSERT. This is necessary because data-modifying statements are only allowed
in WITH clauses that are attached to the top-level statement. However, normal WITH visibility rules
apply, so it is possible to refer to the WITH statement’s output from the sub-SELECT.

Data-modifying statements in WITH usually have RETURNING clauses (see Section 6.4), as shown
in the example above. It is the output of the RETURNING clause, not the target table of the data-modi-
fying statement, that forms the temporary table that can be referred to by the rest of the query. If a
data-modifying statement in WITH lacks a RETURNING clause, then it forms no temporary table and
cannot be referred to in the rest of the query. Such a statement will be executed nonetheless. A not-
particularly-useful example is:

WITH t AS (
    DELETE FROM foo
) DELETE FROM bar;

This example would remove all rows from tables foo and bar. The number of affected rows reported
to the client would only include rows removed from bar.

Recursive self-references in data-modifying statements are not allowed. In some cases it is possible
to work around this limitation by referring to the output of a recursive WITH, for example:

WITH RECURSIVE included_parts(sub_part, part) AS (
    SELECT sub_part, part FROM parts WHERE part = 'our_product'
    UNION ALL
    SELECT p.sub_part, p.part
    FROM included_parts pr, parts p
    WHERE p.part = pr.sub_part
) DELETE FROM parts
    WHERE part IN (SELECT part FROM included_parts);

This query would remove all direct and indirect subparts of a product.

Data-modifying statements in WITH are executed exactly once, and always to completion, indepen-
dently of whether the primary query reads all (or indeed any) of their output. Notice that this is differ-
et from the rule for SELECT in WITH: as stated in the previous section, execution of a SELECT is
carried only as far as the primary query demands its output.

The sub-statements in WITH are executed concurrently with each other and with the main query.
Therefore, when using data-modifying statements in WITH, the order in which the specified updates
actually happen is unpredictable. All the statements are executed with the same snapshot (see Chap-
ter 13), so they cannot “see” one another’s effects on the target tables. This alleviates the effects of the
unpredictability of the actual order of row updates, and means that RETURNING data is the only way
to communicate changes between different WITH sub-statements and the main query. An example of
this is that in

WITH t AS (  UPDATE products SET price = price * 1.05
      RETURNING *
) SELECT * FROM products;

the outer SELECT would return the original prices before the action of the UPDATE, while in
WITH t AS (  
    UPDATE products SET price = price * 1.05  
    RETURNING *
  )
SELECT * FROM t;

the outer SELECT would return the updated data.

Trying to update the same row twice in a single statement is not supported. Only one of the modifications takes place, but it is not easy (and sometimes not possible) to reliably predict which one. This also applies to deleting a row that was already updated in the same statement: only the update is performed. Therefore you should generally avoid trying to modify a single row twice in a single statement. In particular avoid writing WITH sub-statements that could affect the same rows changed by the main statement or a sibling sub-statement. The effects of such a statement will not be predictable.

At present, any table used as the target of a data-modifying statement in WITH must not have a conditional rule, nor an ALSO rule, nor an INSTEAD rule that expands to multiple statements.
Chapter 8. Data Types

PostgreSQL has a rich set of native data types available to users. Users can add new types to PostgreSQL using the CREATE TYPE command.

Table 8.1 shows all the built-in general-purpose data types. Most of the alternative names listed in the “Aliases” column are the names used internally by PostgreSQL for historical reasons. In addition, some internally used or deprecated types are available, but are not listed here.

<table>
<thead>
<tr>
<th>Name</th>
<th>Aliases</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bigint</td>
<td>int8</td>
<td>signed eight-byte integer</td>
</tr>
<tr>
<td>bigserial</td>
<td>serial8</td>
<td>autoincrementing eight-byte integer</td>
</tr>
<tr>
<td>bit [ (n) ]</td>
<td></td>
<td>fixed-length bit string</td>
</tr>
<tr>
<td>bit varying</td>
<td>varbit</td>
<td>variable-length bit string</td>
</tr>
<tr>
<td>bit varying</td>
<td>varbit</td>
<td>variable-length bit string</td>
</tr>
<tr>
<td>boolean</td>
<td>bool</td>
<td>logical Boolean (true/false)</td>
</tr>
<tr>
<td>box</td>
<td></td>
<td>rectangular box on a plane</td>
</tr>
<tr>
<td>bytea</td>
<td></td>
<td>binary data (“byte array”)</td>
</tr>
<tr>
<td>character</td>
<td>char</td>
<td>fixed-length character string</td>
</tr>
<tr>
<td>character</td>
<td>varchar</td>
<td>variable-length character string</td>
</tr>
<tr>
<td>cidr</td>
<td></td>
<td>IPv4 or IPv6 network address</td>
</tr>
<tr>
<td>circle</td>
<td></td>
<td>circle on a plane</td>
</tr>
<tr>
<td>date</td>
<td></td>
<td>calendar date (year, month, day)</td>
</tr>
<tr>
<td>double precision</td>
<td>float8</td>
<td>double precision floating-point number (8 bytes)</td>
</tr>
<tr>
<td>inet</td>
<td></td>
<td>IPv4 or IPv6 host address</td>
</tr>
<tr>
<td>integer</td>
<td>int, int4</td>
<td>signed four-byte integer</td>
</tr>
<tr>
<td>interval</td>
<td></td>
<td>time span</td>
</tr>
<tr>
<td>json</td>
<td></td>
<td>textual JSON data</td>
</tr>
<tr>
<td>jsonb</td>
<td></td>
<td>binary JSON data, decomposed</td>
</tr>
<tr>
<td>line</td>
<td></td>
<td>infinite line on a plane</td>
</tr>
<tr>
<td>lseg</td>
<td></td>
<td>line segment on a plane</td>
</tr>
<tr>
<td>macaddr</td>
<td></td>
<td>MAC (Media Access Control) address</td>
</tr>
<tr>
<td>macaddr8</td>
<td></td>
<td>MAC (Media Access Control) address (EUI-64 format)</td>
</tr>
<tr>
<td>money</td>
<td></td>
<td>currency amount</td>
</tr>
<tr>
<td>numeric</td>
<td>decimal</td>
<td>exact numeric of selectable precision</td>
</tr>
<tr>
<td>path</td>
<td></td>
<td>geometric path on a plane</td>
</tr>
<tr>
<td>pg_lsn</td>
<td></td>
<td>PostgreSQL Log Sequence Number</td>
</tr>
<tr>
<td>point</td>
<td></td>
<td>geometric point on a plane</td>
</tr>
</tbody>
</table>
Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Aliases</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>polygon</td>
<td></td>
<td>closed geometric path on a plane</td>
</tr>
<tr>
<td>real</td>
<td>float4</td>
<td>single precision floating-point number (4 bytes)</td>
</tr>
<tr>
<td>smallint</td>
<td>int2</td>
<td>signed two-byte integer</td>
</tr>
<tr>
<td>smallserial</td>
<td>serial2</td>
<td>autoincrementing two-byte integer</td>
</tr>
<tr>
<td>serial</td>
<td>serial4</td>
<td>autoincrementing four-byte integer</td>
</tr>
<tr>
<td>text</td>
<td></td>
<td>variable-length character string</td>
</tr>
<tr>
<td>time [ (p) ] [ without time zone ]</td>
<td></td>
<td>time of day (no time zone)</td>
</tr>
<tr>
<td>time [ (p) ] with time zone</td>
<td>timetz</td>
<td>time of day, including time zone</td>
</tr>
<tr>
<td>timestamp [ (p) ] [ without time zone ]</td>
<td></td>
<td>date and time (no time zone)</td>
</tr>
<tr>
<td>timestamp [ (p) ] with time zone</td>
<td>timestamptz</td>
<td>date and time, including time zone</td>
</tr>
<tr>
<td>tsquery</td>
<td></td>
<td>text search query</td>
</tr>
<tr>
<td>tsvector</td>
<td></td>
<td>text search document</td>
</tr>
<tr>
<td>txid_snapshot</td>
<td></td>
<td>user-level transaction ID snapshot</td>
</tr>
<tr>
<td>uuid</td>
<td></td>
<td>universally unique identifier</td>
</tr>
<tr>
<td>xml</td>
<td></td>
<td>XML data</td>
</tr>
</tbody>
</table>

**Compatibility**

The following types (or spellings thereof) are specified by SQL: `bigint`, `bit`, `bit varying`, `boolean`, `char`, `character varying`, `character`, `varchar`, `date`, `double precision`, `integer`, `interval`, `numeric`, `decimal`, `real`, `smallint`, `time` (with or without time zone), `timestamp` (with or without time zone), `xml`.

Each data type has an external representation determined by its input and output functions. Many of the built-in types have obvious external formats. However, several types are either unique to PostgreSQL, such as geometric paths, or have several possible formats, such as the date and time types. Some of the input and output functions are not invertible, i.e., the result of an output function might lose accuracy when compared to the original input.

### 8.1. Numeric Types

Numeric types consist of two-, four-, and eight-byte integers, four- and eight-byte floating-point numbers, and selectable-precision decimals. Table 8.2 lists the available types.

**Table 8.2. Numeric Types**

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>smallint</td>
<td>2 bytes</td>
<td>small-range integer</td>
<td>-32768 to +32767</td>
</tr>
</tbody>
</table>
### Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>integer</td>
<td>4 bytes</td>
<td>typical choice for integer</td>
<td>-2147483648 to +2147483647</td>
</tr>
<tr>
<td>bigint</td>
<td>8 bytes</td>
<td>large-range integer</td>
<td>-9223372036854775808 to +9223372036854775807</td>
</tr>
<tr>
<td>decimal</td>
<td>variable</td>
<td>user-specified precision, exact</td>
<td>up to 131072 digits before the decimal point; up to 16383 digits after the decimal point</td>
</tr>
<tr>
<td>numeric</td>
<td>variable</td>
<td>user-specified precision, exact</td>
<td>up to 131072 digits before the decimal point; up to 16383 digits after the decimal point</td>
</tr>
<tr>
<td>real</td>
<td>4 bytes</td>
<td>variable-precision, inexact</td>
<td>6 decimal digits precision</td>
</tr>
<tr>
<td>double precision</td>
<td>8 bytes</td>
<td>variable-precision, inexact</td>
<td>15 decimal digits precision</td>
</tr>
<tr>
<td>smallserial</td>
<td>2 bytes</td>
<td>small autoincrementing integer</td>
<td>1 to 32767</td>
</tr>
<tr>
<td>serial</td>
<td>4 bytes</td>
<td>autoincrementing integer</td>
<td>1 to 2147483647</td>
</tr>
<tr>
<td>bigserial</td>
<td>8 bytes</td>
<td>large autoincrementing integer</td>
<td>9223372036854775807</td>
</tr>
</tbody>
</table>

The syntax of constants for the numeric types is described in Section 4.1.2. The numeric types have a full set of corresponding arithmetic operators and functions. Refer to Chapter 9 for more information. The following sections describe the types in detail.

### 8.1.1. Integer Types

The types smallint, integer, and bigint store whole numbers, that is, numbers without fractional components, of various ranges. Attempts to store values outside of the allowed range will result in an error.

The type integer is the common choice, as it offers the best balance between range, storage size, and performance. The smallint type is generally only used if disk space is at a premium. The bigint type is designed to be used when the range of the integer type is insufficient.

SQL only specifies the integer types integer (or int), smallint, and bigint. The type names int2, int4, and int8 are extensions, which are also used by some other SQL database systems.

### 8.1.2. Arbitrary Precision Numbers

The type numeric can store numbers with a very large number of digits. It is especially recommended for storing monetary amounts and other quantities where exactness is required. Calculations with numeric values yield exact results where possible, e.g., addition, subtraction, multiplication. However, calculations on numeric values are very slow compared to the integer types, or to the floating-point types described in the next section.

We use the following terms below: The precision of a numeric is the total count of significant digits in the whole number, that is, the number of digits to both sides of the decimal point. The scale of a numeric is the count of decimal digits in the fractional part, to the right of the decimal point. So the number 23.5141 has a precision of 6 and a scale of 4. Integers can be considered to have a scale of zero.
Both the maximum precision and the maximum scale of a numeric column can be configured. To declare a column of type numeric use the syntax:

\[
\text{NUMERIC(\text{precision}, \text{scale})}
\]

The precision must be positive, the scale zero or positive. Alternatively:

\[
\text{NUMERIC(\text{precision})}
\]

selects a scale of 0. Specifying:

\[
\text{NUMERIC}
\]

without any precision or scale creates a column in which numeric values of any precision and scale can be stored, up to the implementation limit on precision. A column of this kind will not coerce input values to any particular scale, whereas numeric columns with a declared scale will coerce input values to that scale. (The SQL standard requires a default scale of 0, i.e., coercion to integer precision. We find this a bit useless. If you're concerned about portability, always specify the precision and scale explicitly.)

**Note**

The maximum allowed precision when explicitly specified in the type declaration is 1000; \text{NUMERIC} without a specified precision is subject to the limits described in Table 8.2.

If the scale of a value to be stored is greater than the declared scale of the column, the system will round the value to the specified number of fractional digits. Then, if the number of digits to the left of the decimal point exceeds the declared precision minus the declared scale, an error is raised.

Numeric values are physically stored without any extra leading or trailing zeroes. Thus, the declared precision and scale of a column are maximums, not fixed allocations. (In this sense the numeric type is more akin to \text{varchar(\text{n})} than to \text{char(\text{n})}. The actual storage requirement is two bytes for each group of four decimal digits, plus three to eight bytes overhead.

In addition to ordinary numeric values, the numeric type allows the special value NaN, meaning “not-a-number”. Any operation on NaN yields another NaN. When writing this value as a constant in an SQL command, you must put quotes around it, for example \text{UPDATE table SET x = 'NaN'}. On input, the string NaN is recognized in a case-insensitive manner.

**Note**

In most implementations of the “not-a-number” concept, NaN is not considered equal to any other numeric value (including NaN). In order to allow numeric values to be sorted and used in tree-based indexes, PostgreSQL treats NaN values as equal, and greater than all non-NaN values.

The types \text{decimal} and \text{numeric} are equivalent. Both types are part of the SQL standard.

When rounding values, the numeric type rounds ties away from zero, while (on most machines) the real and double precision types round ties to the nearest even number. For example:
SELECT x,
    round(x::numeric) AS num_round,
    round(x::double precision) AS dbl_round
FROM generate_series(-3.5, 3.5, 1) as x;

<table>
<thead>
<tr>
<th></th>
<th>num_round</th>
<th>dbl_round</th>
</tr>
</thead>
<tbody>
<tr>
<td>-3.5</td>
<td>-4</td>
<td>-4</td>
</tr>
<tr>
<td>-2.5</td>
<td>-3</td>
<td>-2</td>
</tr>
<tr>
<td>-1.5</td>
<td>-2</td>
<td>-2</td>
</tr>
<tr>
<td>-0.5</td>
<td>-1</td>
<td>-0</td>
</tr>
<tr>
<td>0.5</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>1.5</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>2.5</td>
<td>3</td>
<td>2</td>
</tr>
<tr>
<td>3.5</td>
<td>4</td>
<td>4</td>
</tr>
</tbody>
</table>

(8 rows)

8.1.3. Floating-Point Types

The data types real and double precision are inexact, variable-precision numeric types. In practice, these types are usually implementations of IEEE Standard 754 for Binary Floating-Point Arithmetic (single and double precision, respectively), to the extent that the underlying processor, operating system, and compiler support it.

Inexact means that some values cannot be converted exactly to the internal format and are stored as approximations, so that storing and retrieving a value might show slight discrepancies. Managing these errors and how they propagate through calculations is the subject of an entire branch of mathematics and computer science and will not be discussed here, except for the following points:

- If you require exact storage and calculations (such as for monetary amounts), use the numeric type instead.
- If you want to do complicated calculations with these types for anything important, especially if you rely on certain behavior in boundary cases (infinity, underflow), you should evaluate the implementation carefully.
- Comparing two floating-point values for equality might not always work as expected.

On most platforms, the real type has a range of at least 1E-37 to 1E+37 with a precision of at least 6 decimal digits. The double precision type typically has a range of around 1E-307 to 1E+308 with a precision of at least 15 digits. Values that are too large or too small will cause an error. Rounding might take place if the precision of an input number is too high. Numbers too close to zero that are not representable as distinct from zero will cause an underflow error.

Note

The extra_float_digits setting controls the number of extra significant digits included when a floating point value is converted to text for output. With the default value of 0, the output is the same on every platform supported by PostgreSQL. Increasing it will produce output that more accurately represents the stored value, but may be unportable.

In addition to ordinary numeric values, the floating-point types have several special values:

Infinity
-Infinity
NaN

These represent the IEEE 754 special values “infinity”, “negative infinity”, and “not-a-number”, respectively. (On a machine whose floating-point arithmetic does not follow IEEE 754, these values will probably not work as expected.) When writing these values as constants in an SQL command, you must put quotes around them, for example UPDATE table SET x = '-Infinity'. On input, these strings are recognized in a case-insensitive manner.

Note

IEEE754 specifies that NaN should not compare equal to any other floating-point value (including NaN). In order to allow floating-point values to be sorted and used in tree-based indexes, PostgreSQL treats NaN values as equal, and greater than all non-NaN values.

PostgreSQL also supports the SQL-standard notations float and float(p) for specifying inexact numeric types. Here, p specifies the minimum acceptable precision in binary digits. PostgreSQL accepts float(1) to float(24) as selecting the real type, while float(25) to float(53) select double precision. Values of p outside the allowed range draw an error. float with no precision specified is taken to mean double precision.

Note

The assumption that real and double precision have exactly 24 and 53 bits in the mantissa respectively is correct for IEEE-standard floating point implementations. On non-IEEE platforms it might be off a little, but for simplicity the same ranges of p are used on all platforms.

8.1.4. Serial Types

Note

This section describes a PostgreSQL-specific way to create an autoincrementing column. Another way is to use the SQL-standard identity column feature, described at CREATE TABLE.

The data types smallserial, serial and bigserial are not true types, but merely a notational convenience for creating unique identifier columns (similar to the AUTO_INCREMENT property supported by some other databases). In the current implementation, specifying:

CREATE TABLE tablename (  
  colname SERIAL  
);

is equivalent to specifying:

CREATE SEQUENCE tablename_colname_seq;  
CREATE TABLE tablename (  
  colname integer NOT NULL DEFAULT  
  nextval('tablename_colname_seq')  
);
Thus, we have created an integer column and arranged for its default values to be assigned from a sequence generator. A NOT NULL constraint is applied to ensure that a null value cannot be inserted. (In most cases you would also want to attach a UNIQUE or PRIMARY KEY constraint to prevent duplicate values from being inserted by accident, but this is not automatic.) Lastly, the sequence is marked as "owned by" the column, so that it will be dropped if the column or table is dropped.

Note

Because smallserial, serial and bigserial are implemented using sequences, there may be "holes" or gaps in the sequence of values which appears in the column, even if no rows are ever deleted. A value allocated from the sequence is still "used up" even if a row containing that value is never successfully inserted into the table column. This may happen, for example, if the inserting transaction rolls back. See nextval() in Section 9.16 for details.

To insert the next value of the sequence into the serial column, specify that the serial column should be assigned its default value. This can be done either by excluding the column from the list of columns in the INSERT statement, or through the use of the DEFAULT key word.

The type names serial and serial4 are equivalent: both create integer columns. The type names bigserial and serial8 work the same way, except that they create a bigint column. bigserial should be used if you anticipate the use of more than $2^{31}$ identifiers over the lifetime of the table. The type names smallserial and serial2 also work the same way, except that they create a smallint column.

The sequence created for a serial column is automatically dropped when the owning column is dropped. You can drop the sequence without dropping the column, but this will force removal of the column default expression.

8.2. Monetary Types

The money type stores a currency amount with a fixed fractional precision; see Table 8.3. The fractional precision is determined by the database's lc_monetary setting. The range shown in the table assumes there are two fractional digits. Input is accepted in a variety of formats, including integer and floating-point literals, as well as typical currency formatting, such as '$1,000.00'. Output is generally in the latter form but depends on the locale.

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>money</td>
<td>8 bytes</td>
<td>currency amount</td>
<td>-92233720368547758.08 to +92233720368547758.07</td>
</tr>
</tbody>
</table>

Since the output of this data type is locale-sensitive, it might not work to load money data into a database that has a different setting of lc_monetary. To avoid problems, before restoring a dump into a new database make sure lc_monetary has the same or equivalent value as in the database that was dumped.

Values of the numeric, int, and bigint data types can be cast to money. Conversion from the real and double precision data types can be done by casting to numeric first, for example:
Data Types

```
SELECT '12.34'::float8::numeric::money;
```

However, this is not recommended. Floating point numbers should not be used to handle money due to the potential for rounding errors.

A `money` value can be cast to `numeric` without loss of precision. Conversion to other types could potentially lose precision, and must also be done in two stages:

```
SELECT '52093.89'::money::numeric::float8;
```

Division of a `money` value by an integer value is performed with truncation of the fractional part towards zero. To get a rounded result, divide by a floating-point value, or cast the `money` value to `numeric` before dividing and back to `money` afterwards. (The latter is preferable to avoid risking precision loss.) When a `money` value is divided by another `money` value, the result is double precision (i.e., a pure number, not money); the currency units cancel each other out in the division.

8.3. Character Types

Table 8.4. Character Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character varying(n), varchar(n)</td>
<td>variable-length with limit</td>
</tr>
<tr>
<td>character(n), char(n)</td>
<td>fixed-length, blank padded</td>
</tr>
<tr>
<td>text</td>
<td>variable unlimited length</td>
</tr>
</tbody>
</table>

Table 8.4 shows the general-purpose character types available in PostgreSQL.

SQL defines two primary character types: `character varying(n)` and `character(n)`, where `n` is a positive integer. Both of these types can store strings up to `n` characters (not bytes) in length. An attempt to store a longer string into a column of these types will result in an error, unless the excess characters are all spaces, in which case the string will be truncated to the maximum length. (This somewhat bizarre exception is required by the SQL standard.) If the string to be stored is shorter than the declared length, values of type `character` will be space-padded; values of type `character varying` will simply store the shorter string.

If one explicitly casts a value to `character varying(n)` or `character(n)`, then an over-length value will be truncated to `n` characters without raising an error. (This too is required by the SQL standard.)

The notations `varchar(n)` and `char(n)` are aliases for `character varying(n)` and `character(n)`, respectively. `character` without length specifier is equivalent to `character(1)`. If `character varying` is used without length specifier, the type accepts strings of any size. The latter is a PostgreSQL extension.

In addition, PostgreSQL provides the `text` type, which stores strings of any length. Although the type `text` is not in the SQL standard, several other SQL database management systems have it as well.

Values of type `character` are physically padded with spaces to the specified width `n`, and are stored and displayed that way. However, trailing spaces are treated as semantically insignificant and disregarded when comparing two values of type `character`. In collations where whitespace is significant, this behavior can produce unexpected results; for example, `SELECT 'a'::CHAR(2) collate "C" < E'a
'::CHAR(2)` returns true, even though `C` locale would consider a space to be greater than a newline. Trailing spaces are removed when converting a `character` value to one of the other string types. Note that trailing spaces are semantically significant in `character varying` and `text` values, and when using pattern matching, that is `LIKE` and regular expressions.
The storage requirement for a short string (up to 126 bytes) is 1 byte plus the actual string, which includes the space padding in the case of character. Longer strings have 4 bytes of overhead instead of 1. Long strings are compressed by the system automatically, so the physical requirement on disk might be less. Very long values are also stored in background tables so that they do not interfere with rapid access to shorter column values. In any case, the longest possible character string that can be stored is about 1 GB. (The maximum value that will be allowed for $n$ in the data type declaration is less than that. It wouldn't be useful to change this because with multibyte character encodings the number of characters and bytes can be quite different. If you desire to store long strings with no specific upper limit, use text or character varying without a length specifier, rather than making up an arbitrary length limit.)

**Tip**

There is no performance difference among these three types, apart from increased storage space when using the blank-padded type, and a few extra CPU cycles to check the length when storing into a length-constrained column. While character($n$) has performance advantages in some other database systems, there is no such advantage in PostgreSQL; in fact character($n$) is usually the slowest of the three because of its additional storage costs. In most situations text or character varying should be used instead.

Refer to Section 4.1.2.1 for information about the syntax of string literals, and to Chapter 9 for information about available operators and functions. The database character set determines the character set used to store textual values; for more information on character set support, refer to Section 23.3.

**Example 8.1. Using the Character Types**

```sql
CREATE TABLE test1 (a character(4));
INSERT INTO test1 VALUES ('ok');
SELECT a, char_length(a) FROM test1; --

<table>
<thead>
<tr>
<th>a</th>
<th>char_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>ok</td>
<td>2</td>
</tr>
</tbody>
</table>
```

```sql
CREATE TABLE test2 (b varchar(5));
INSERT INTO test2 VALUES ('ok');
INSERT INTO test2 VALUES ('good      ');
INSERT INTO test2 VALUES ('too long');
ERROR: value too long for type character varying(5)
INSERT INTO test2 VALUES ('too long':varchar(5)); -- explicit truncation
SELECT b, char_length(b) FROM test2;

<table>
<thead>
<tr>
<th>b</th>
<th>char_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>ok</td>
<td>2</td>
</tr>
<tr>
<td>good</td>
<td>5</td>
</tr>
<tr>
<td>too l</td>
<td>5</td>
</tr>
</tbody>
</table>
```

The `char_length` function is discussed in Section 9.4.

There are two other fixed-length character types in PostgreSQL, shown in Table 8.5. The `name` type exists only for the storage of identifiers in the internal system catalogs and is not intended for use...
by the general user. Its length is currently defined as 64 bytes (63 usable characters plus terminator) but should be referenced using the constant NAMEDATALEN in C source code. The length is set at compile time (and is therefore adjustable for special uses); the default maximum length might change in a future release. The type "char" (note the quotes) is different from char(1) in that it only uses one byte of storage. It is internally used in the system catalogs as a simplistic enumeration type.

Table 8.5. Special Character Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&quot;char&quot;</td>
<td>1 byte</td>
<td>single-byte internal type</td>
</tr>
<tr>
<td>name</td>
<td>64 bytes</td>
<td>internal type for object names</td>
</tr>
</tbody>
</table>

8.4. Binary Data Types

The bytea data type allows storage of binary strings; see Table 8.6.

Table 8.6. Binary Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytea</td>
<td>1 or 4 bytes plus the actual binary string</td>
<td>variable-length binary string</td>
</tr>
</tbody>
</table>

A binary string is a sequence of octets (or bytes). Binary strings are distinguished from character strings in two ways. First, binary strings specifically allow storing octets of value zero and other "non-printable" octets (usually, octets outside the decimal range 32 to 126). Character strings disallow zero octets, and also disallow any other octet values and sequences of octet values that are invalid according to the database's selected character set encoding. Second, operations on binary strings process the actual bytes, whereas the processing of character strings depends on locale settings. In short, binary strings are appropriate for storing data that the programmer thinks of as "raw bytes", whereas character strings are appropriate for storing text.

The bytea type supports two formats for input and output: “hex” format and PostgreSQL's historical “escape” format. Both of these are always accepted on input. The output format depends on the configuration parameter bytea_output; the default is hex. (Note that the hex format was introduced in PostgreSQL 9.0; earlier versions and some tools don’t understand it.)

The SQL standard defines a different binary string type, called BLOB or BINARY LARGE OBJECT. The input format is different from bytea, but the provided functions and operators are mostly the same.

8.4.1. bytea Hex Format

The “hex” format encodes binary data as 2 hexadecimal digits per byte, most significant nibble first. The entire string is preceded by the sequence \\x (to distinguish it from the escape format). In some contexts, the initial backslash may need to be escaped by doubling it, in the same cases in which backslashes have to be doubled in escape format; details appear below. The hexadecimal digits can be either upper or lower case, and whitespace is permitted between digit pairs (but not within a digit pair nor in the starting \\x sequence). The hex format is compatible with a wide range of external applications and protocols, and it tends to be faster to convert than the escape format, so its use is preferred.

Example:

SELECT '\xDEADBEEF';
8.4.2. bytea Escape Format

The “escape” format is the traditional PostgreSQL format for the bytea type. It takes the approach of representing a binary string as a sequence of ASCII characters, while converting those bytes that cannot be represented as an ASCII character into special escape sequences. If, from the point of view of the application, representing bytes as characters makes sense, then this representation can be convenient. But in practice it is usually confusing because it fuzzes up the distinction between binary strings and character strings, and also the particular escape mechanism that was chosen is somewhat unwieldy. Therefore, this format should probably be avoided for most new applications.

When entering bytea values in escape format, octets of certain values must be escaped, while all octet values can be escaped. In general, to escape an octet, convert it into its three-digit octal value and precede it by a backslash (or two backslashes, if writing the value as a literal using escape string syntax). Backslash itself (octet decimal value 92) can alternatively be represented by double backslashes. Table 8.7 shows the characters that must be escaped, and gives the alternative escape sequences where applicable.

<table>
<thead>
<tr>
<th>Decimal Value</th>
<th>Octet Description</th>
<th>Escaped Representation</th>
<th>Input Example</th>
<th>Output Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>zero octet</td>
<td>\000</td>
<td>SELECT \000::bytea;</td>
<td>\x00</td>
</tr>
<tr>
<td>39</td>
<td>single quote</td>
<td>''' or '\047'</td>
<td>SELECT '''::bytea;</td>
<td>\x27</td>
</tr>
<tr>
<td>92</td>
<td>backslash</td>
<td>'\ or '\134'</td>
<td>SELECT '::bytea;</td>
<td>\x5c</td>
</tr>
<tr>
<td>0 to 31 and 127 to “non-printable” octets</td>
<td>\xxx (octal value)</td>
<td>(octal select value)</td>
<td>(octal select value)</td>
<td>(octal select value)</td>
</tr>
<tr>
<td>255</td>
<td></td>
<td>\001</td>
<td>\001::bytea;</td>
<td>\x01</td>
</tr>
</tbody>
</table>

The requirement to escape non-printable octets varies depending on locale settings. In some instances you can get away with leaving them unescaped. Note that the result in each of the examples in Table 8.7 was exactly one octet in length, even though the output representation is sometimes more than one character.

The reason multiple backslashes are required, as shown in Table 8.7, is that an input string written as a string literal must pass through two parse phases in the PostgreSQL server. The first backslash of each pair is interpreted as an escape character by the string-literal parser (assuming escape string syntax is used) and is therefore consumed, leaving the second backslash of the pair. (Dollar-quoted strings can be used to avoid this level of escaping.) The remaining backslash is then recognized by the bytea input function as starting either a three digit octal value or escaping another backslash. For example, a string literal passed to the server as \001 becomes \001 after passing through the escape string parser. The \001 is then sent to the bytea input function, where it is converted to a single octet with a decimal value of 1. Note that the single-quote character is not treated specially by bytea, so it follows the normal rules for string literals. (See also Section 4.1.2.1.)

Bytea octets are output in hex format by default. If you change bytea_output to escape, “non-printable” octet are converted to equivalent three-digit octal value and preceded by one backslash. Most “printable” octets are output by their standard representation in the client character set, e.g.:

```
SET bytea_output = 'escape';

SELECT 'abc \153\154\155 \052\251\124'::bytea;
bytea
----------------
141
```
The octet with decimal value 92 (backslash) is doubled in the output. Details are in Table 8.8.

Table 8.8. bytea Output Escaped Octets

<table>
<thead>
<tr>
<th>Decimal Value</th>
<th>Octet Description</th>
<th>Escaped Output Representation</th>
<th>Output Example</th>
<th>Output Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>92</td>
<td>backslash</td>
<td><code>\</code></td>
<td>SELECT <code>\134</code>::bytea;</td>
<td></td>
</tr>
<tr>
<td>0 to 31 and 127 to 255</td>
<td>“non-printable” octets</td>
<td><code>\xxx</code> (octal value)</td>
<td>SELECT <code>\001</code>::bytea;</td>
<td></td>
</tr>
<tr>
<td>32 to 126</td>
<td>“printable” octets</td>
<td>client character set representation</td>
<td>SELECT <code>\176</code>::bytea;</td>
<td></td>
</tr>
</tbody>
</table>

Depending on the front end to PostgreSQL you use, you might have additional work to do in terms of escaping and unescaping bytea strings. For example, you might also have to escape line feeds and carriage returns if your interface automatically translates these.

8.5. Date/Time Types

PostgreSQL supports the full set of SQL date and time types, shown in Table 8.9. The operations available on these data types are described in Section 9.9. Dates are counted according to the Gregorian calendar, even in years before that calendar was introduced (see Section B.4 for more information).

Table 8.9. Date/Time Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Low Value</th>
<th>High Value</th>
<th>Resolution</th>
</tr>
</thead>
<tbody>
<tr>
<td>timestamp</td>
<td>8 bytes</td>
<td>both date and time (no time zone)</td>
<td>4713 BC</td>
<td>294276 AD</td>
<td>1 microsecond</td>
</tr>
<tr>
<td>[ (p) ]</td>
<td></td>
<td>with time zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>timestamp</td>
<td>8 bytes</td>
<td>both date and time, with time zone</td>
<td>4713 BC</td>
<td>294276 AD</td>
<td>1 microsecond</td>
</tr>
<tr>
<td>[ (p) ]</td>
<td></td>
<td>with time zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>date</td>
<td>4 bytes</td>
<td>date (no time of day)</td>
<td>4713 BC</td>
<td>5874897 AD</td>
<td>1 day</td>
</tr>
<tr>
<td>time</td>
<td>8 bytes</td>
<td>time of day (no date)</td>
<td>00:00:00</td>
<td>24:00:00</td>
<td>1 microsecond</td>
</tr>
<tr>
<td>[ (p) ]</td>
<td></td>
<td>with time zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>time</td>
<td>12 bytes</td>
<td>time of day (no date), with time zone</td>
<td>00:00:00+1459</td>
<td>24:00:00-1459</td>
<td>1 microsecond</td>
</tr>
<tr>
<td>[ (p) ]</td>
<td></td>
<td>with time zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>interval</td>
<td>16 bytes</td>
<td>time interval</td>
<td>-1780000000 years</td>
<td>178000000 years</td>
<td>1 microsecond</td>
</tr>
<tr>
<td>[ fields ]</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>[ (p) ]</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Note

The SQL standard requires that writing just timestamp be equivalent to timestamp without time zone, and PostgreSQL honors that behavior. time-
Data Types

stamptz is accepted as an abbreviation for timestamp with time zone; this is a PostgreSQL extension.

time, timestamp, and interval accept an optional precision value $p$ which specifies the number of fractional digits retained in the seconds field. By default, there is no explicit bound on precision. The allowed range of $p$ is from 0 to 6.

The interval type has an additional option, which is to restrict the set of stored fields by writing one of these phrases:

YEAR
MONTH
DAY
HOUR
MINUTE
SECOND
YEAR TO MONTH
DAY TO HOUR
DAY TO MINUTE
DAY TO SECOND
HOUR TO MINUTE
HOUR TO SECOND
MINUTE TO SECOND

Note that if both fields and $p$ are specified, the fields must include SECOND, since the precision applies only to the seconds.

The type time with time zone is defined by the SQL standard, but the definition exhibits properties which lead to questionable usefulness. In most cases, a combination of date, time, timestamp without time zone, and timestamp with time zone should provide a complete range of date/time functionality required by any application.

The types abstime and reltime are lower precision types which are used internally. You are discouraged from using these types in applications; these internal types might disappear in a future release.

8.5.1. Date/Time Input

Date and time input is accepted in almost any reasonable format, including ISO 8601, SQL-compatible, traditional POSTGRES, and others. For some formats, ordering of day, month, and year in date input is ambiguous and there is support for specifying the expected ordering of these fields. Set the DateStyle parameter to MDY to select month-day-year interpretation, DMY to select day-month-year interpretation, or YMD to select year-month-day interpretation.

PostgreSQL is more flexible in handling date/time input than the SQL standard requires. See Appendix B for the exact parsing rules of date/time input and for the recognized text fields including months, days of the week, and time zones.

Remember that any date or time literal input needs to be enclosed in single quotes, like text strings. Refer to Section 4.1.2.7 for more information. SQL requires the following syntax

\[ \text{type [ (p) ] \ 'value'} \]

where $p$ is an optional precision specification giving the number of fractional digits in the seconds field. Precision can be specified for time, timestamp, and interval types, and can range from 0 to 6. If no precision is specified in a constant specification, it defaults to the precision of the literal value (but not more than 6 digits).
8.5.1.1. Dates

Table 8.10 shows some possible inputs for the date type.

**Table 8.10. Date Input**

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1999-01-08</td>
<td>ISO 8601; January 8 in any mode (recommended format)</td>
</tr>
<tr>
<td>January 8, 1999</td>
<td>unambiguous in any datestyle input mode</td>
</tr>
<tr>
<td>1/8/1999</td>
<td>January 8 in MDY mode; August 1 in DMY mode</td>
</tr>
<tr>
<td>1/18/1999</td>
<td>January 18 in MDY mode; rejected in other modes</td>
</tr>
<tr>
<td>01/02/03</td>
<td>January 2, 2003 in MDY mode; February 1, 2003 in DMY mode; February 3, 2001 in YMD mode</td>
</tr>
<tr>
<td>1999-Jan-08</td>
<td>January 8 in any mode</td>
</tr>
<tr>
<td>Jan-08-1999</td>
<td>January 8 in any mode</td>
</tr>
<tr>
<td>08-Jan-1999</td>
<td>January 8 in any mode</td>
</tr>
<tr>
<td>99-Jan-08</td>
<td>January 8 in YMD mode, else error</td>
</tr>
<tr>
<td>08-Jan-99</td>
<td>January 8, except error in YMD mode</td>
</tr>
<tr>
<td>Jan-08-99</td>
<td>January 8, except error in YMD mode</td>
</tr>
<tr>
<td>19990108</td>
<td>ISO 8601; January 8, 1999 in any mode</td>
</tr>
<tr>
<td>990108</td>
<td>ISO 8601; January 8, 1999 in any mode</td>
</tr>
<tr>
<td>1999.008</td>
<td>year and day of year</td>
</tr>
<tr>
<td>J2451187</td>
<td>Julian date</td>
</tr>
<tr>
<td>January 8, 99 BC</td>
<td>year 99 BC</td>
</tr>
</tbody>
</table>

8.5.1.2. Times

The time-of-day types are `time [ (p) ]` without time zone and `time [ (p) ]` with time zone. Time alone is equivalent to time without time zone.

Valid input for these types consists of a time of day followed by an optional time zone. (See Table 8.11 and Table 8.12.) If a time zone is specified in the input for `time without time zone`, it is silently ignored. You can also specify a date but it will be ignored, except when you use a time zone name that involves a daylight-savings rule, such as America/New_York. In this case specifying the date is required in order to determine whether standard or daylight-savings time applies. The appropriate time zone offset is recorded in the `time with time zone` value.

**Table 8.11. Time Input**

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>04:05:06.789</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>040506</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05 AM</td>
<td>same as 04:05; AM does not affect value</td>
</tr>
<tr>
<td>04:05 PM</td>
<td>same as 16:05; input hour must be &lt;= 12</td>
</tr>
<tr>
<td>04:05:06.789-8</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06-08:00</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>Example</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------</td>
<td>-------------------------------------------------------</td>
</tr>
<tr>
<td>04:05-08:00</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>040506-08</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06 PST</td>
<td>time zone specified by abbreviation</td>
</tr>
<tr>
<td>2003-04-12 04:05:06 America/New_York</td>
<td>time zone specified by full name</td>
</tr>
</tbody>
</table>

### Table 8.12. Time Zone Input

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PST</td>
<td>Abbreviation (for Pacific Standard Time)</td>
</tr>
<tr>
<td>America/New_York</td>
<td>Full time zone name</td>
</tr>
<tr>
<td>PST8PDT</td>
<td>POSIX-style time zone specification</td>
</tr>
<tr>
<td>-8:00</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>-800</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>-8</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>zulu</td>
<td>Military abbreviation for UTC</td>
</tr>
<tr>
<td>z</td>
<td>Short form of zulu</td>
</tr>
</tbody>
</table>

Refer to Section 8.5.3 for more information on how to specify time zones.

### 8.5.1.3. Time Stamps

Valid input for the time stamp types consists of the concatenation of a date and a time, followed by an optional time zone, followed by an optional AD or BC. (Alternatively, AD/BC can appear before the time zone, but this is not the preferred ordering.) Thus:

1999-01-08 04:05:06

and:

1999-01-08 04:05:06 -8:00

are valid values, which follow the ISO 8601 standard. In addition, the common format:

January 8 04:05:06 1999 PST

is supported.

The SQL standard differentiates `timestamp without time zone` and `timestamp with time zone` literals by the presence of a “+” or “-” symbol and time zone offset after the time. Hence, according to the standard,

`TIMESTAMP '2004-10-19 10:23:54'` is a `timestamp without time zone`, while

`TIMESTAMP '2004-10-19 10:23:54+02'` is a `timestamp with time zone`. PostgreSQL never examines the content of a literal string before determining its type, and therefore will treat both of the above as `timestamp without time zone`. To ensure that a literal is treated as `timestamp with time zone`, give it the correct explicit type:
TIMESTAMP WITH TIME ZONE '2004-10-19 10:23:54+02'

In a literal that has been determined to be timestamp without time zone, PostgreSQL will silently ignore any time zone indication. That is, the resulting value is derived from the date/time fields in the input value, and is not adjusted for time zone.

For timestamp with time zone, the internally stored value is always in UTC (Universal Coordinated Time, traditionally known as Greenwich Mean Time, GMT). An input value that has an explicit time zone specified is converted to UTC using the appropriate offset for that time zone. If no time zone is stated in the input string, then it is assumed to be in the time zone indicated by the system’s TimeZone parameter, and is converted to UTC using the offset for the timezone zone.

When a timestamp with time zone value is output, it is always converted from UTC to the current timezone zone, and displayed as local time in that zone. To see the time in another time zone, either change timezone or use the AT TIME ZONE construct (see Section 9.9.3).

Conversions between timestamp without time zone and timestamp with time zone normally assume that the timestamp without time zone value should be taken or given as timezone local time. A different time zone can be specified for the conversion using AT TIME ZONE.

8.5.1.4. Special Values

PostgreSQL supports several special date/time input values for convenience, as shown in Table 8.13. The values infinity and -infinity are specially represented inside the system and will be displayed unchanged; but the others are simply notational shorthands that will be converted to ordinary date/time values when read. (In particular, now and related strings are converted to a specific time value as soon as they are read.) All of these values need to be enclosed in single quotes when used as constants in SQL commands.

Table 8.13. Special Date/Time Inputs

<table>
<thead>
<tr>
<th>Input String</th>
<th>Valid Types</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>epoch</td>
<td>date, timestamp</td>
<td>1970-01-01 00:00:00+00 (Unix system time zero)</td>
</tr>
<tr>
<td>infinity</td>
<td>date, timestamp</td>
<td>later than all other time stamps</td>
</tr>
<tr>
<td>-infinity</td>
<td>date, timestamp</td>
<td>earlier than all other time stamps</td>
</tr>
<tr>
<td>now</td>
<td>date, time, timestamp</td>
<td>current transaction’s start time</td>
</tr>
<tr>
<td>today</td>
<td>date, timestamp</td>
<td>midnight today</td>
</tr>
<tr>
<td>tomorrow</td>
<td>date, timestamp</td>
<td>midnight tomorrow</td>
</tr>
<tr>
<td>yesterday</td>
<td>date, timestamp</td>
<td>midnight yesterday</td>
</tr>
<tr>
<td>allballs</td>
<td>time</td>
<td>00:00:00.00 UTC</td>
</tr>
</tbody>
</table>

The following SQL-compatible functions can also be used to obtain the current time value for the corresponding data type: CURRENT_DATE, CURRENT_TIME, CURRENT_TIMESTAMP, LOCALTIME, LOCALTIMESTAMP. The latter four accept an optional subsecond precision specification. (See Section 9.9.4.) Note that these are SQL functions and are not recognized in data input strings.

8.5.2. Date/Time Output

The output format of the date/time types can be set to one of the four styles ISO 8601, SQL (Ingres), traditional POSTGRES (Unix date format), or German. The default is the ISO format. (The SQL standard requires the use of the ISO 8601 format. The name of the “SQL” output format is a historical accident.) Table 8.14 shows examples of each output style. The output of the date and time types is
generally only the date or time part in accordance with the given examples. However, the POSTGRES style outputs date-only values in ISO format.

### Table 8.14. Date/Time Output Styles

<table>
<thead>
<tr>
<th>Style Specification</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>ISO</td>
<td>ISO 8601, SQL standard</td>
<td>1997-12-17 07:37:16-08</td>
</tr>
<tr>
<td>SQL</td>
<td>traditional style</td>
<td>12/17/1997 07:37:16.00 PST</td>
</tr>
<tr>
<td>Postgres</td>
<td>original style</td>
<td>Wed Dec 17 07:37:16 1997 PST</td>
</tr>
<tr>
<td>German</td>
<td>regional style</td>
<td>17.12.1997 07:37:16.00 PST</td>
</tr>
</tbody>
</table>

**Note**

ISO 8601 specifies the use of uppercase letter T to separate the date and time. PostgreSQL accepts that format on input, but on output it uses a space rather than T, as shown above. This is for readability and for consistency with RFC 3339 as well as some other database systems.

In the SQL and POSTGRES styles, day appears before month if DMY field ordering has been specified, otherwise month appears before day. (See Section 8.5.1 for how this setting also affects interpretation of input values.) Table 8.15 shows examples.

### Table 8.15. Date Order Conventions

<table>
<thead>
<tr>
<th>datestyle Setting</th>
<th>Input Ordering</th>
<th>Example Output</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL, DMY</td>
<td>day/month/year</td>
<td>17/12/1997 07:37:16.00 CET</td>
</tr>
<tr>
<td>SQL, MDY</td>
<td>month/day/year</td>
<td>12/17/1997 07:37:16.00 PST</td>
</tr>
<tr>
<td>Postgres, DMY</td>
<td>day/month/year</td>
<td>Wed 17 Dec 07:37:16 1997 PST</td>
</tr>
</tbody>
</table>

The date/time style can be selected by the user using the SET datestyle command, the DateStyle parameter in the postgresql.conf configuration file, or the PGDATESTYLE environment variable on the server or client.

The formatting function `to_char` (see Section 9.8) is also available as a more flexible way to format date/time output.

### 8.5.3. Time Zones

Time zones, and time-zone conventions, are influenced by political decisions, not just earth geometry. Time zones around the world became somewhat standardized during the 1900s, but continue to be prone to arbitrary changes, particularly with respect to daylight-savings rules. PostgreSQL uses the widely-used IANA (Olson) time zone database for information about historical time zone rules. For times in the future, the assumption is that the latest known rules for a given time zone will continue to be observed indefinitely far into the future.

PostgreSQL endeavors to be compatible with the SQL standard definitions for typical usage. However, the SQL standard has an odd mix of date and time types and capabilities. Two obvious problems are:
Although the date type cannot have an associated time zone, the time type can. Time zones in the real world have little meaning unless associated with a date as well as a time, since the offset can vary through the year with daylight-saving time boundaries.

The default time zone is specified as a constant numeric offset from UTC. It is therefore impossible to adapt to daylight-saving time when doing date/time arithmetic across DST boundaries.

To address these difficulties, we recommend using date/time types that contain both date and time when using time zones. We do not recommend using the type time with time zone (though it is supported by PostgreSQL for legacy applications and for compliance with the SQL standard). PostgreSQL assumes your local time zone for any type containing only date or time.

All timezone-aware dates and times are stored internally in UTC. They are converted to local time in the zone specified by the TimeZone configuration parameter before being displayed to the client.

PostgreSQL allows you to specify time zones in three different forms:

- A full time zone name, for example America/New_York. The recognized time zone names are listed in the pg_timezone_names view (see Section 52.90). PostgreSQL uses the widely-used IANA time zone data for this purpose, so the same time zone names are also recognized by much other software.

- A time zone abbreviation, for example PST. Such a specification merely defines a particular offset from UTC, in contrast to full time zone names which can imply a set of daylight savings transition-date rules as well. The recognized abbreviations are listed in the pg_timezone_abbrs view (see Section 52.89). You cannot set the configuration parameters TimeZone or log_timezone to a time zone abbreviation, but you can use abbreviations in date/time input values and with the AT TIME ZONE operator.

- In addition to the timezone names and abbreviations, PostgreSQL will accept POSIX-style time zone specifications of the form STD offset or STD offset DST, where STD is a zone abbreviation, offset is a numeric offset in hours west from UTC, and DST is an optional daylight-savings zone abbreviation, assumed to stand for one hour ahead of the given offset. For example, if EST5EDT were not already a recognized zone name, it would be accepted and would be functionally equivalent to United States East Coast time. In this syntax, a zone abbreviation can be a string of letters, or an arbitrary string surrounded by angle brackets (<>). When a daylight-savings zone abbreviation is present, it is assumed to be used according to the same daylight-savings transition rules used in the IANA time zone database's posixrules entry. In a standard PostgreSQL installation, posixrules is the same as US/Eastern, so that POSIX-style time zone specifications follow USA daylight-savings rules. If needed, you can adjust this behavior by replacing the posixrules file.

In short, this is the difference between abbreviations and full names: abbreviations represent a specific offset from UTC, whereas many of the full names imply a local daylight-savings time rule, and so have two possible UTC offsets. As an example, 2014-06-04 12:00 America/New_York represents noon local time in New York, which for this particular date was Eastern Daylight Time (UTC-4). So 2014-06-04 12:00 EDT specifies that same time instant. But 2014-06-04 12:00 EST specifies noon Eastern Standard Time (UTC-5), regardless of whether daylight savings was nominally in effect on that date.

To complicate matters, some jurisdictions have used the same timezone abbreviation to mean different UTC offsets at different times; for example, in Moscow MSK has meant UTC+3 in some years and UTC+4 in others. PostgreSQL interprets such abbreviations according to whatever they meant (or had most recently meant) on the specified date; but, as with the EST example above, this is not necessarily the same as local civil time on that date.

One should be wary that the POSIX-style time zone feature can lead to silently accepting bogus input, since there is no check on the reasonableness of the zone abbreviations. For example, SET TIME-ZONE TO FOOBAR0 will work, leaving the system effectively using a rather peculiar abbreviation for UTC. Another issue to keep in mind is that in POSIX time zone names, positive offsets are used...
for locations west of Greenwich. Everywhere else, PostgreSQL follows the ISO-8601 convention that positive timezone offsets are east of Greenwich.

In all cases, timezone names and abbreviations are recognized case-insensitively. (This is a change from PostgreSQL versions prior to 8.2, which were case-sensitive in some contexts but not others.)

Neither timezone names nor abbreviations are hard-wired into the server; they are obtained from configuration files stored under `.../share/timezone/` and `.../share/timezonesets/` of the installation directory (see Section B.3).

The TimeZone configuration parameter can be set in the file `postgresql.conf`, or in any of the other standard ways described in Chapter 19. There are also some special ways to set it:

- The SQL command `SET TIME ZONE` sets the time zone for the session. This is an alternative spelling of `SET TIMEZONE TO` with a more SQL-spec-compatible syntax.
- The `PGTZ` environment variable is used by libpq clients to send a `SET TIME ZONE` command to the server upon connection.

### 8.5.4. Interval Input

Interval values can be written using the following verbose syntax:

```
[@] quantity unit [quantity unit...] [direction]
```

where `quantity` is a number (possibly signed); `unit` is `microsecond`, `millisecond`, `second`, `minute`, `hour`, `day`, `week`, `month`, `year`, `decade`, `century`, `millennium`, or abbreviations or plurals of these units; `direction` can be `ago` or empty. The at sign (`@`) is optional noise. The amounts of the different units are implicitly added with appropriate sign accounting. `ago` negates all the fields. This syntax is also used for interval output, if `IntervalStyle` is set to `postgres_verbose`.

Quantities of days, hours, minutes, and seconds can be specified without explicit unit markings. For example, `'1 12:59:10'` is read the same as `'1 day 12 hours 59 min 10 sec'`. Also, a combination of years and months can be specified with a dash; for example `'200-10'` is read the same as `'200 years 10 months'`. (These shorter forms are in fact the only ones allowed by the SQL standard, and are used for output when `IntervalStyle` is set to `sql_standard`.)

Interval values can also be written as ISO 8601 time intervals, using either the “format with designators” of the standard’s section 4.4.3.2 or the “alternative format” of section 4.4.3.3. The format with designators looks like this:

```
P quantity unit [ quantity unit ...] [ T [ quantity unit ...]]
```

The string must start with a `P`, and may include a `T` that introduces the time-of-day units. The available unit abbreviations are given in Table 8.16. Units may be omitted, and may be specified in any order, but units smaller than a day must appear after `T`. In particular, the meaning of `M` depends on whether it is before or after `T`.

### Table 8.16. ISO 8601 Interval Unit Abbreviations

<table>
<thead>
<tr>
<th>Abbreviation</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>Y</td>
<td>Years</td>
</tr>
<tr>
<td>M</td>
<td>Months (in the date part)</td>
</tr>
<tr>
<td>W</td>
<td>Weeks</td>
</tr>
<tr>
<td>D</td>
<td>Days</td>
</tr>
</tbody>
</table>

---

149
<table>
<thead>
<tr>
<th>Abbreviation</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>H</td>
<td>Hours</td>
</tr>
<tr>
<td>M</td>
<td>Minutes (in the time part)</td>
</tr>
<tr>
<td>S</td>
<td>Seconds</td>
</tr>
</tbody>
</table>

In the alternative format:

\[ P \ [ \text{years-months-days} ] \ [ T \text{hours:minutes:seconds} ] \]

the string must begin with \( P \), and a \( T \) separates the date and time parts of the interval. The values are given as numbers similar to ISO 8601 dates.

When writing an interval constant with a \( \text{fields} \) specification, or when assigning a string to an interval column that was defined with a \( \text{fields} \) specification, the interpretation of unmarked quantities depends on the \( \text{fields} \). For example, \texttt{INTERVAL '1' YEAR} is read as 1 year, whereas \texttt{INTERVAL '1' means 1 second. Also, field values “to the right” of the least significant field allowed by the \( \text{fields} \) specification are silently discarded. For example, writing \texttt{INTERVAL '1 day 2:03:04' HOUR TO MINUTE} results in dropping the seconds field, but not the day field.

According to the SQL standard all fields of an interval value must have the same sign, so a leading negative sign applies to all fields; for example the negative sign in the interval literal \(' -1 2:03:04' \) applies to both the days and hour/minute/second parts. PostgreSQL allows the fields to have different signs, and traditionally treats each field in the textual representation as independently signed, so that the hour/minute/second part is considered positive in this example. If \( \text{IntervalStyle} \) is set to \( \text{sql_standard} \) then a leading sign is considered to apply to all fields (but only if no additional signs appear). Otherwise the traditional PostgreSQL interpretation is used. To avoid ambiguity, it’s recommended to attach an explicit sign to each field if any field is negative.

In the verbose input format, and in some fields of the more compact input formats, field values can have fractional parts; for example \('1.5 \text{ week}' or '01:02:03.45'\). Such input is converted to the appropriate number of months, days, and seconds for storage. When this would result in a fractional number of months or days, the fraction is added to the lower-order fields using the conversion factors 1 month = 30 days and 1 day = 24 hours. For example, \('1.5 \text{ month}' becomes 1 month and 15 days. Only seconds will ever be shown as fractional on output.

Table 8.17 shows some examples of valid interval input.

**Table 8.17. Interval Input**

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1-2</td>
<td>SQL standard format: 1 year 2 months</td>
</tr>
<tr>
<td>3 4:05:06</td>
<td>SQL standard format: 3 days 4 hours 5 minutes 6 seconds</td>
</tr>
<tr>
<td>1 year 2 months 3 days 4 hours 5 minutes 6 seconds</td>
<td>Traditional Postgres format: 1 year 2 months 3 days 4 hours 5 minutes 6 seconds</td>
</tr>
<tr>
<td>P1Y2M3DT4H15M6S</td>
<td>ISO 8601 “format with designators”: same meaning as above</td>
</tr>
<tr>
<td>P0001-02-03T04:05:06</td>
<td>ISO 8601 “alternative format”: same meaning as above</td>
</tr>
</tbody>
</table>

Internally interval values are stored as months, days, and seconds. This is done because the number of days in a month varies, and a day can have 23 or 25 hours if a daylight savings time adjustment is involved. The months and days fields are integers while the seconds field can store fractions. Because intervals are usually created from constant strings or \( \text{timestamp} \) subtraction, this storage method works well in most cases, but can cause unexpected results:
SELECT EXTRACT(hours from '80 minutes'::interval);
  date_part
  -----------
  1

SELECT EXTRACT(days from '80 hours'::interval);
  date_part
  -----------
  0

Functions justify_days and justify_hours are available for adjusting days and hours that
overflow their normal ranges.

8.5.5. Interval Output

The output format of the interval type can be set to one of the four styles sql_standard, post-
gres, postgres_verbose, or iso_8601, using the command SET intervalstyle. The
default is the postgres format. Table 8.18 shows examples of each output style.

The sql_standard style produces output that conforms to the SQL standard's specification for
interval literal strings, if the interval value meets the standard's restrictions (either year-month only or
day-time only, with no mixing of positive and negative components). Otherwise the output looks like
a standard year-month literal string followed by a day-time literal string, with explicit signs added to
disambiguate mixed-sign intervals.

The output of the postgres style matches the output of PostgreSQL releases prior to 8.4 when the
DateStyle parameter was set to ISO.

The output of the postgres_verbose style matches the output of PostgreSQL releases prior to
8.4 when the DateStyle parameter was set to non-ISO output.

The output of the iso_8601 style matches the “format with designators” described in section 4.4.3.2
of the ISO 8601 standard.

Table 8.18. Interval Output Style Examples

<table>
<thead>
<tr>
<th>Style Specification</th>
<th>Year-Month Interval</th>
<th>Day-Time Interval</th>
<th>Mixed Interval</th>
</tr>
</thead>
<tbody>
<tr>
<td>sql_standard</td>
<td>1-2</td>
<td>3 days 04:05:06</td>
<td>-1 year -2 mons +3 days-04:05:06</td>
</tr>
<tr>
<td>postgres</td>
<td>1 year 2 mons</td>
<td>3 days 04:05:06</td>
<td>-1 year -2 mons +3 days-04:05:06</td>
</tr>
<tr>
<td>postgres_verbose</td>
<td>@ 1 year 2 mons</td>
<td>@ 3 days 4 hours 5 mins 6 secs</td>
<td>@ 1 year 2 mons -3 days 4 hours 5 mins 6 secs ago</td>
</tr>
<tr>
<td>iso_8601</td>
<td>P1Y2M</td>
<td>P3DT4H5M6S</td>
<td>P-1Y-2M3DT-4H-5M-6S</td>
</tr>
</tbody>
</table>

8.6. Boolean Type

PostgreSQL provides the standard SQL type boolean; see Table 8.19. The boolean type can have
several states: “true”, “false”, and a third state, “unknown”, which is represented by the SQL null value.

Table 8.19. Boolean Data Type

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>boolean</td>
<td>1 byte</td>
<td>state of true or false</td>
</tr>
</tbody>
</table>

Valid literal values for the “true” state are:
For the “false” state, the following values can be used:

FALSE
'f'
'false'
'n'
'no'
'off'
'0'

Leading or trailing whitespace is ignored, and case does not matter. The key words TRUE and FALSE are the preferred (SQL-compliant) usage.

Example 8.2 shows that boolean values are output using the letters t and f.

Example 8.2. Using the boolean Type

```
CREATE TABLE test1 (a boolean, b text);
INSERT INTO test1 VALUES (TRUE, 'sic est');
INSERT INTO test1 VALUES (FALSE, 'non est');
SELECT * FROM test1;
  a |    b
----+---------
  t | sic est
  f | non est

SELECT * FROM test1 WHERE a;
  a |    b
----+---------
  t | sic est
```

8.7. Enumerated Types

Enumerated (enum) types are data types that comprise a static, ordered set of values. They are equivalent to the enum types supported in a number of programming languages. An example of an enum type might be the days of the week, or a set of status values for a piece of data.

8.7.1. Declaration of Enumerated Types

Enum types are created using the CREATE TYPE command, for example:

```
CREATE TYPE mood AS ENUM ('sad', 'ok', 'happy');
```

Once created, the enum type can be used in table and function definitions much like any other type:

```
CREATE TYPE mood AS ENUM ('sad', 'ok', 'happy');
CREATE TABLE person (
    name text,
```
8.7.2. Ordering

The ordering of the values in an enum type is the order in which the values were listed when the type was created. All standard comparison operators and related aggregate functions are supported for enums. For example:

```sql
INSERT INTO person VALUES ('Larry', 'sad');
INSERT INTO person VALUES ('Curly', 'ok');
SELECT * FROM person WHERE current_mood > 'sad';
```

<table>
<thead>
<tr>
<th>name</th>
<th>current_mood</th>
</tr>
</thead>
<tbody>
<tr>
<td>Moe</td>
<td>happy</td>
</tr>
<tr>
<td>Curly</td>
<td>ok</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
(2 rows)

```sql
SELECT * FROM person WHERE current_mood > 'sad' ORDER BY current_mood;
```

<table>
<thead>
<tr>
<th>name</th>
<th>current_mood</th>
</tr>
</thead>
<tbody>
<tr>
<td>Curly</td>
<td>ok</td>
</tr>
<tr>
<td>Moe</td>
<td>happy</td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
(2 rows)

```sql
SELECT name FROM person WHERE current_mood = (SELECT MIN(current_mood) FROM person);
```

<table>
<thead>
<tr>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Larry</td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
</tbody>
</table>
(1 row)

8.7.3. Type Safety

Each enumerated data type is separate and cannot be compared with other enumerated types. See this example:

```sql
CREATE TYPE happiness AS ENUM ('happy', 'very happy', 'ecstatic');
CREATE TABLE holidays (num_weeks integer, happiness happiness);
INSERT INTO holidays(num_weeks,happiness) VALUES (4, 'happy');
INSERT INTO holidays(num_weeks,happiness) VALUES (6, 'very happy');
INSERT INTO holidays(num_weeks,happiness) VALUES (8, 'ecstatic');
INSERT INTO holidays(num_weeks,happiness) VALUES (2, 'sad');
ERROR: invalid input value for enum happiness: "sad"
```

```sql
SELECT person.name, holidays.num_weeks FROM person, holidays
```
WHERE person.current_mood = holidays.happiness;
ERROR: operator does not exist: mood = happiness

If you really need to do something like that, you can either write a custom operator or add explicit casts to your query:

SELECT person.name, holidays.num_weeks FROM person, holidays
WHERE person.current_mood::text = holidays.happiness::text;

<table>
<thead>
<tr>
<th>name</th>
<th>num_weeks</th>
</tr>
</thead>
<tbody>
<tr>
<td>Moe</td>
<td>4</td>
</tr>
</tbody>
</table>

(1 row)

8.7.4. Implementation Details

Enum labels are case sensitive, so 'happy' is not the same as 'HAPPY'. White space in the labels is significant too.

Although enum types are primarily intended for static sets of values, there is support for adding new values to an existing enum type, and for renaming values (see ALTER TYPE). Existing values cannot be removed from an enum type, nor can the sort ordering of such values be changed, short of dropping and re-creating the enum type.

An enum value occupies four bytes on disk. The length of an enum value's textual label is limited by the NAMEDATALEN setting compiled into PostgreSQL; in standard builds this means at most 63 bytes.

The translations from internal enum values to textual labels are kept in the system catalog pg_enum. Querying this catalog directly can be useful.

8.8. Geometric Types

Geometric data types represent two-dimensional spatial objects. Table 8.20 shows the geometric types available in PostgreSQL.

### Table 8.20. Geometric Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>point</td>
<td>16 bytes</td>
<td>Point on a plane</td>
<td>(x,y)</td>
</tr>
<tr>
<td>line</td>
<td>32 bytes</td>
<td>Infinite line</td>
<td>{A,B,C}</td>
</tr>
<tr>
<td>lseg</td>
<td>32 bytes</td>
<td>Finite line segment</td>
<td>((x1,y1),(x2,y2))</td>
</tr>
<tr>
<td>box</td>
<td>32 bytes</td>
<td>Rectangular box</td>
<td>((x1,y1),(x2,y2))</td>
</tr>
<tr>
<td>path</td>
<td>16+16n bytes</td>
<td>Closed path (similar to ((x1,y1),...)) polygon</td>
<td></td>
</tr>
<tr>
<td>path</td>
<td>16+16n bytes</td>
<td>Open path</td>
<td>[(x1,y1),...]</td>
</tr>
<tr>
<td>polygon</td>
<td>40+16n bytes</td>
<td>Polygon (similar to ((x1,y1),...)) closed path</td>
<td></td>
</tr>
<tr>
<td>circle</td>
<td>24 bytes</td>
<td>Circle</td>
<td>&lt;(x,y),r&gt; (center point and radius)</td>
</tr>
</tbody>
</table>

A rich set of functions and operators is available to perform various geometric operations such as scaling, translation, rotation, and determining intersections. They are explained in Section 9.11.

8.8.1. Points
Points are the fundamental two-dimensional building block for geometric types. Values of type `point` are specified using either of the following syntaxes:

\[(x, y)\]

where `x` and `y` are the respective coordinates, as floating-point numbers.

Points are output using the first syntax.

### 8.8.2. Lines

Lines are represented by the linear equation `Ax + By + C = 0`, where `A` and `B` are not both zero. Values of type `line` are input and output in the following form:

\[\{ A, B, C \}\]

Alternatively, any of the following forms can be used for input:

\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[x_1, y_1, x_2, y_2\]\]

where \((x_1, y_1)\) and \((x_2, y_2)\) are two different points on the line.

### 8.8.3. Line Segments

Line segments are represented by pairs of points that are the endpoints of the segment. Values of type `lseg` are specified using any of the following syntaxes:

\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[x_1, y_1, x_2, y_2\]\]

where \((x_1, y_1)\) and \((x_2, y_2)\) are the end points of the line segment.

Line segments are output using the first syntax.

### 8.8.4. Boxes

Boxes are represented by pairs of points that are opposite corners of the box. Values of type `box` are specified using any of the following syntaxes:

\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[(x_1, y_1), (x_2, y_2)\]\]
\[\[x_1, y_1, x_2, y_2\]\]

where \((x_1, y_1)\) and \((x_2, y_2)\) are any two opposite corners of the box.

Boxes are output using the second syntax.

Any two opposite corners can be supplied on input, but the values will be reordered as needed to store the upper right and lower left corners, in that order.
8.8.5. Paths

Paths are represented by lists of connected points. Paths can be open, where the first and last points in the list are considered not connected, or closed, where the first and last points are considered connected.

Values of type `path` are specified using any of the following syntaxes:

\[
[ (x_1, y_1), \ldots, (x_n, y_n) ]
\]
\[
( (x_1, y_1), \ldots, (x_n, y_n) )
\]
\[
( x_1, y_1, \ldots, x_n, y_n )
\]
\[
( x_1, y_1, \ldots, x_n, y_n )
\]

where the points are the end points of the line segments comprising the path. Square brackets `[]` indicate an open path, while parentheses `()` indicate a closed path. When the outermost parentheses are omitted, as in the third through fifth syntaxes, a closed path is assumed.

Paths are output using the first or second syntax, as appropriate.

8.8.6. Polygons

Polygons are represented by lists of points (the vertexes of the polygon). Polygons are very similar to closed paths, but are stored differently and have their own set of support routines.

Values of type `polygon` are specified using any of the following syntaxes:

\[
( (x_1, y_1), \ldots, (x_n, y_n) )
\]
\[
( x_1, y_1, \ldots, x_n, y_n )
\]
\[
( x_1, y_1, \ldots, x_n, y_n )
\]
\[
x_1, y_1, \ldots, x_n, y_n
\]

where the points are the end points of the line segments comprising the boundary of the polygon.

Polygons are output using the first syntax.

8.8.7. Circles

Circles are represented by a center point and radius. Values of type `circle` are specified using any of the following syntaxes:

\[
< (x, y), r >
\]
\[
( (x, y), r )
\]
\[
( x, y ), r
\]
\[
x, y , r
\]

where \((x, y)\) is the center point and \(r\) is the radius of the circle.

Circles are output using the first syntax.

8.9. Network Address Types

PostgreSQL offers data types to store IPv4, IPv6, and MAC addresses, as shown in Table 8.21. It is better to use these types instead of plain text types to store network addresses, because these types offer input error checking and specialized operators and functions (see Section 9.12).
Table 8.21. Network Address Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>cidr</td>
<td>7 or 19 bytes</td>
<td>IPv4 and IPv6 networks</td>
</tr>
<tr>
<td>inet</td>
<td>7 or 19 bytes</td>
<td>IPv4 and IPv6 hosts and networks</td>
</tr>
<tr>
<td>macaddr</td>
<td>6 bytes</td>
<td>MAC addresses</td>
</tr>
<tr>
<td>macaddr8</td>
<td>8 bytes</td>
<td>MAC addresses (EUI-64 format)</td>
</tr>
</tbody>
</table>

When sorting inet or cidr data types, IPv4 addresses will always sort before IPv6 addresses, including IPv4 addresses encapsulated or mapped to IPv6 addresses, such as ::10.2.3.4 or ::ffff:10.4.3.2.

### 8.9.1. inet

The inet type holds an IPv4 or IPv6 host address, and optionally its subnet, all in one field. The subnet is represented by the number of network address bits present in the host address (the “netmask”). If the netmask is 32 and the address is IPv4, then the value does not indicate a subnet, only a single host. In IPv6, the address length is 128 bits, so 128 bits specify a unique host address. Note that if you want to accept only networks, you should use the cidr type rather than inet.

The input format for this type is \textit{address/y} where \textit{address} is an IPv4 or IPv6 address and \textit{y} is the number of bits in the netmask. If the /y portion is missing, the netmask is 32 for IPv4 and 128 for IPv6, so the value represents just a single host. On display, the /y portion is suppressed if the netmask specifies a single host.

### 8.9.2. cidr

The cidr type holds an IPv4 or IPv6 network specification. Input and output formats follow Classless Internet Domain Routing conventions. The format for specifying networks is \textit{address/y} where \textit{address} is the network represented as an IPv4 or IPv6 address, and \textit{y} is the number of bits in the netmask. If \textit{y} is omitted, it is calculated using assumptions from the older classful network numbering system, except it will be at least large enough to include all of the octets written in the input. It is an error to specify a network address that has bits set to the right of the specified netmask.

Table 8.22 shows some examples.

### Table 8.22. cidr Type Input Examples

<table>
<thead>
<tr>
<th>cidr Input</th>
<th>cidr Output</th>
<th>abbrev(cidr)</th>
</tr>
</thead>
<tbody>
<tr>
<td>192.168.100.128/25</td>
<td>192.168.100.128/25</td>
<td>192.168.100.128/25</td>
</tr>
<tr>
<td>192.168/24</td>
<td>192.168.0/24</td>
<td>192.168.0/24</td>
</tr>
<tr>
<td>192.168/25</td>
<td>192.168.0/25</td>
<td>192.168.0/25</td>
</tr>
<tr>
<td>192.168.1</td>
<td>192.168.1/24</td>
<td>192.168.1/24</td>
</tr>
<tr>
<td>192.168</td>
<td>192.168.0/24</td>
<td>192.168.0/24</td>
</tr>
<tr>
<td>128.1</td>
<td>128.1/16</td>
<td>128.1/16</td>
</tr>
<tr>
<td>128</td>
<td>128.0/16</td>
<td>128.0/16</td>
</tr>
<tr>
<td>128.1.2</td>
<td>128.1.2/24</td>
<td>128.1.2/24</td>
</tr>
<tr>
<td>10.1.2</td>
<td>10.1.2/24</td>
<td>10.1.2/24</td>
</tr>
<tr>
<td>10.1</td>
<td>10.1/16</td>
<td>10.1/16</td>
</tr>
<tr>
<td>10</td>
<td>10.0/8</td>
<td>10/8</td>
</tr>
<tr>
<td>10.1.2.3/32</td>
<td>10.1.2.3/32</td>
<td>10.1.2.3/32</td>
</tr>
<tr>
<td>2001:4f8:3:ba::/64</td>
<td>2001:4f8:3:ba::/64</td>
<td>2001:4f8:3:ba::/64</td>
</tr>
</tbody>
</table>
Data Types

<table>
<thead>
<tr>
<th>cidr Input</th>
<th>cidr Output</th>
<th>abbrev(cidr)</th>
</tr>
</thead>
<tbody>
<tr>
<td>::ffff:1.2.3.0/120</td>
<td>::ffff:1.2.3.0/120</td>
<td>::ffff:1.2.3.0/120</td>
</tr>
<tr>
<td>::ffff:1.2.3.0/128</td>
<td>::ffff:1.2.3.0/128</td>
<td>::ffff:1.2.3.0/128</td>
</tr>
</tbody>
</table>

8.9.3. *inet vs. cidr*

The essential difference between *inet* and *cidr* data types is that *inet* accepts values with nonzero bits to the right of the netmask, whereas *cidr* does not. For example, 192.168.0.1/24 is valid for *inet* but not for *cidr*.

**Tip**

If you do not like the output format for *inet* or *cidr* values, try the functions *host*, *text*, and *abbrev*.

8.9.4. *macaddr*

The *macaddr* type stores MAC addresses, known for example from Ethernet card hardware addresses (although MAC addresses are used for other purposes as well). Input is accepted in the following formats:

- `'08:00:2b:01:02:03'`
- `'08-00-2b-01-02-03'`
- `'08002b:010203'`
- `'08002b-010203'`
- `'0800.2b01.0203'`
- `'0800-2b01-0203'`
- `'08002b010203'`

These examples would all specify the same address. Upper and lower case is accepted for the digits a through f. Output is always in the first of the forms shown.

IEEE Std 802-2001 specifies the second shown form (with hyphens) as the canonical form for MAC addresses, and specifies the first form (with colons) as the bit-reversed notation, so that 08-00-2b-01-02-03 = 01:00:4D:08:04:0C. This convention is widely ignored nowadays, and it is relevant only for obsolete network protocols (such as Token Ring). PostgreSQL makes no provisions for bit reversal, and all accepted formats use the canonical LSB order.

The remaining five input formats are not part of any standard.

8.9.5. *macaddr8*

The *macaddr8* type stores MAC addresses in EUI-64 format, known for example from Ethernet card hardware addresses (although MAC addresses are used for other purposes as well). This type can accept both 6 and 8 byte length MAC addresses and stores them in 8 byte length format. MAC addresses given in 6 byte format will be stored in 8 byte length format with the 4th and 5th bytes set to FF and FE, respectively. Note that IPv6 uses a modified EUI-64 format where the 7th bit should be set to one after the conversion from EUI-48. The function *macaddr8_set7bit* is provided to make this change. Generally speaking, any input which is comprised of pairs of hex digits (on byte boundaries), optionally separated consistently by one of ' : ', '- ', or '. ', is accepted. The number of hex digits must be either 16 (8 bytes) or 12 (6 bytes). Leading and trailing whitespace is ignored. The following are examples of input formats that are accepted:
Data Types

These examples would all specify the same address. Upper and lower case is accepted for the digits a through f. Output is always in the first of the forms shown. The last six input formats that are mentioned above are not part of any standard. To convert a traditional 48 bit MAC address in EUI-48 format to modified EUI-64 format to be included as the host portion of an IPv6 address, use `macaddr8_set7bit` as shown:

```
SELECT macaddr8_set7bit('08:00:2b:01:02:03');
```

```
macaddr8_set7bit
--------------------
0a:00:2b:ff:fe:01:02:03
```

8.10. Bit String Types

Bit strings are strings of 1’s and 0’s. They can be used to store or visualize bit masks. There are two SQL bit types: `bit(n)` and `bit varying(n)`, where n is a positive integer.

`bit` type data must match the length n exactly; it is an error to attempt to store shorter or longer bit strings. `bit varying` data is of variable length up to the maximum length n; longer strings will be rejected. Writing `bit` without a length is equivalent to `bit(1)`, while `bit varying` without a length specification means unlimited length.

Note

If one explicitly casts a bit-string value to `bit(n)`, it will be truncated or zero-padded on the right to be exactly n bits, without raising an error. Similarly, if one explicitly casts a bit-string value to `bit varying(n)`, it will be truncated on the right if it is more than n bits.

Refer to Section 4.1.2.5 for information about the syntax of bit string constants. Bit-logical operators and string manipulation functions are available; see Section 9.6.

Example 8.3. Using the Bit String Types

```
CREATE TABLE test (a BIT(3), b BIT VARYING(5));
INSERT INTO test VALUES (B'101', B'00');
ERROR:  bit string length 2 does not match type bit(3)
INSERT INTO test VALUES (B'10'::bit(3), B'101');
SELECT * FROM test;
a | b
```
A bit string value requires 1 byte for each group of 8 bits, plus 5 or 8 bytes overhead depending on the length of the string (but long values may be compressed or moved out-of-line, as explained in Section 8.3 for character strings).

### 8.11. Text Search Types

PostgreSQL provides two data types that are designed to support full text search, which is the activity of searching through a collection of natural-language documents to locate those that best match a query. The `tsvector` type represents a document in a form optimized for text search; the `tsquery` type similarly represents a text query. Chapter 12 provides a detailed explanation of this facility, and Section 9.13 summarizes the related functions and operators.

#### 8.11.1. tsvector

A `tsvector` value is a sorted list of distinct lexemes, which are words that have been normalized to merge different variants of the same word (see Chapter 12 for details). Sorting and duplicate-elimination are done automatically during input, as shown in this example:

```sql
SELECT 'a fat cat sat on a mat and ate a fat rat'::tsvector;
```

```sql
tsvector
-----------------------------
'a' 'and' 'ate' 'cat' 'fat' 'mat' 'on' 'rat' 'sat'
```

To represent lexemes containing whitespace or punctuation, surround them with quotes:

```sql
SELECT $$the lexeme '    ' contains spaces$$::tsvector;
```

```sql
tsvector
-------------------------------------------
'    ' 'contains' 'lexeme' 'spaces' 'the'
```

(We use dollar-quoted string literals in this example and the next one to avoid the confusion of having to double quote marks within the literals.) Embedded quotes and backslashes must be doubled:

```sql
SELECT $$the lexeme 'Joe''s' contains a quote$$::tsvector;
```

```sql
tsvector
------------------------------------------------
'Joe''s' 'a' 'contains' 'lexeme' 'quote' 'the'
```

Optionally, integer positions can be attached to lexemes:

```sql
SELECT 'a:1 fat:2 cat:3 sat:4 on:5 a:6 mat:7 and:8 ate:9 a:10 fat:11 rat:12':::tsvector;
```

```sql
tsvector
------------------------------------------------------------------------
'a':1,6,10 'and':8 'ate':9 'cat':3 'fat':2,11 'mat':7 'on':5 'rat':12 'sat':4
```

A position normally indicates the source word's location in the document. Positional information can be used for proximity ranking. Position values can range from 1 to 16383; larger numbers are silently set to 16383. Duplicate positions for the same lexeme are discarded.
Lexemes that have positions can further be labeled with a weight, which can be A, B, C, or D. D is the default and hence is not shown on output:

```
SELECT 'a:1A fat:2B,4C cat:5D'::tsvector;
```

```
  tsvector
---------------------
   'a':1A 'cat':5 'fat':2B,4C
```

Weights are typically used to reflect document structure, for example by marking title words differently from body words. Text search ranking functions can assign different priorities to the different weight markers.

It is important to understand that the tsvector type itself does not perform any word normalization; it assumes the words it is given are normalized appropriately for the application. For example,

```
SELECT 'The Fat Rats'::tsvector;
```

```
  tsvector
------------------
   'Fat' 'Rats' 'The'
```

For most English-text-searching applications the above words would be considered non-normalized, but tsvector doesn't care. Raw document text should usually be passed through to_tsvector to normalize the words appropriately for searching:

```
SELECT to_tsvector('english', 'The Fat Rats');
```

```
  to_tsvector
-----------------
   'fat':2 'rat':3
```

Again, see Chapter 12 for more detail.

### 8.11.2. tsquery

A tsquery value stores lexemes that are to be searched for, and can combine them using the Boolean operators & (AND), | (OR), and ! (NOT), as well as the phrase search operator <-> (FOLLOWED BY). There is also a variant <N> of the FOLLOWED BY operator, where N is an integer constant that specifies the distance between the two lexemes being searched for. <-> is equivalent to <1>.

Parentheses can be used to enforce grouping of these operators. In the absence of parentheses, ! (NOT) binds most tightly, <-> (FOLLOWED BY) next most tightly, then & (AND), with | (OR) binding the least tightly.

Here are some examples:

```
SELECT 'fat & rat'::tsquery;
```

```
  tsquery
----------
   'fat' & 'rat'
```

```
SELECT 'fat & (rat | cat)'::tsquery;
```

```
  tsquery
----------
   'fat' & ('rat' | 'cat')
```

```
SELECT 'fat & rat & ! cat'::tsquery;
```
Data Types

**tsquery**

`'fat' & 'rat' & '!cat'`

Optionally, lexemes in a `tsquery` can be labeled with one or more weight letters, which restricts them to match only `tsvector` lexemes with one of those weights:

```sql
SELECT 'fat:ab & cat':tsquery;

--

'fat':AB & 'cat'
```

Also, lexemes in a `tsquery` can be labeled with `*` to specify prefix matching:

```sql
SELECT 'super:*':tsquery;

--

'super': *
```

This query will match any word in a `tsvector` that begins with “super”.

Quoting rules for lexemes are the same as described previously for lexemes in `tsvector`; and, as with `tsvector`, any required normalization of words must be done before converting to the `tsquery` type. The `to_tsquery` function is convenient for performing such normalization:

```sql
SELECT to_tsquery('Fat:ab & Cats');

--

'fat':AB & 'cat'
```

Note that `to_tsquery` will process prefixes in the same way as other words, which means this comparison returns true:

```sql
SELECT to_tsvector( 'postgraduate' ) @@ to_tsquery( 'postgres:*' );

--

true
```

because `postgres` gets stemmed to `postgr`:

```sql
SELECT to_tsvector( 'postgraduate' ), to_tsquery( 'postgres:*' );

to_tsvector | to_tsquery
-------------------
'postgradu':1 | 'postgr': *
```

which will match the stemmed form of `postgraduate`.

### 8.12. UUID Type

The data type `uuid` stores Universally Unique Identifiers (UUID) as defined by RFC 4122, ISO/IEC 9834-8:2005, and related standards. (Some systems refer to this data type as a globally unique identifier, or GUID, instead.) This identifier is a 128-bit quantity that is generated by an algorithm chosen to make it very unlikely that the same identifier will be generated by anyone else in the known universe using the same algorithm. Therefore, for distributed systems, these identifiers provide a better uniqueness guarantee than sequence generators, which are only unique within a single database.
A UUID is written as a sequence of lower-case hexadecimal digits, in several groups separated by hyphens, specifically a group of 8 digits followed by three groups of 4 digits followed by a group of 12 digits, for a total of 32 digits representing the 128 bits. An example of a UUID in this standard form is:

a0eebc99-9c0b-4ef8-bb6d-6bb9bd380a11

PostgreSQL also accepts the following alternative forms for input: use of upper-case digits, the standard format surrounded by braces, omitting some or all hyphens, adding a hyphen after any group of four digits. Examples are:

A0EEBC99-9C0B-4EF8-BB6D-6BB9BD380A11
{a0eebc99-9c0b-4ef8-bb6d-6bb9bd380a11}
a0ee-bc99-9c0b-4ef8-bb6d-6bb9-bd38-0a11
{a0eebc99-9c0b4ef8-bb6d6bb9-bd380a11}

Output is always in the standard form.

PostgreSQL provides storage and comparison functions for UUIDs, but the core database does not include any function for generating UUIDs, because no single algorithm is well suited for every application. The uuid-ossp module provides functions that implement several standard algorithms. The pgcrypto module also provides a generation function for random UUIDs. Alternatively, UUIDs could be generated by client applications or other libraries invoked through a server-side function.

8.13. XML Type

The xml data type can be used to store XML data. Its advantage over storing XML data in a text field is that it checks the input values for well-formedness, and there are support functions to perform type-safe operations on it; see Section 9.14. Use of this data type requires the installation to have been built with configure --with-libxml.

The xml type can store well-formed “documents”, as defined by the XML standard, as well as “content” fragments, which are defined by the production XMLDecl? content in the XML standard. Roughly, this means that content fragments can have more than one top-level element or character node. The expression xmlvalue IS DOCUMENT can be used to evaluate whether a particular xml value is a full document or only a content fragment.

8.13.1. Creating XML Values

To produce a value of type xml from character data, use the function xmlparse:

XMLPARSE ( ( DOCUMENT | CONTENT ) value)

Examples:

XMLPARSE (DOCUMENT '<?xml version="1.0"?><book><title>Manual</title><chapter>...</chapter></book>')
XMLPARSE (CONTENT '"abc<foo>bar</foo><bar>foo</bar>"')

While this is the only way to convert character strings into XML values according to the SQL standard, the PostgreSQL-specific syntaxes:

xml '<foo>bar</foo>'
'<?foo>bar</foo>'::xml
can also be used.

The xml type does not validate input values against a document type declaration (DTD), even when
the input value specifies a DTD. There is also currently no built-in support for validating against other
XML schema languages such as XML Schema.

The inverse operation, producing a character string value from xml, uses the function xmlserial-
ize:

\[
\text{XMLSERIALIZE} \left( \left\{ \text{DOCUMENT} \mid \text{CONTENT} \right\} \text{value} \mid \text{AS} \mid \text{type} \right)
\]

`type` can be `character`, `character varying`, or `text` (or an alias for one of those). Again, according to the SQL standard, this is the only way to convert between type xml and character types,
but PostgreSQL also allows you to simply cast the value.

When a character string value is cast to or from type xml without going through XMLPARSE or XM-
LSERIALIZE, respectively, the choice of `DOCUMENT` versus `CONTENT` is determined by the “XML
option” session configuration parameter, which can be set using the standard command:

\[
\text{SET XML OPTION} \left( \left\{ \text{DOCUMENT} \mid \text{CONTENT} \right\} \right); \\
\text{or the more PostgreSQL-like syntax}
\]

\[
\text{SET xmloption TO} \left( \left\{ \text{DOCUMENT} \mid \text{CONTENT} \right\} \right);
\]

The default is `CONTENT`, so all forms of XML data are allowed.

**Note**

With the default XML option setting, you cannot directly cast character strings to type
xml if they contain a document type declaration, because the definition of XML con-
tent fragment does not accept them. If you need to do that, either use XMLPARSE or
change the XML option.

### 8.13.2. Encoding Handling

Care must be taken when dealing with multiple character encodings on the client, server, and in the
XML data passed through them. When using the text mode to pass queries to the server and query
results to the client (which is the normal mode), PostgreSQL converts all character data passed be-
tween the client and the server and vice versa to the character encoding of the respective end; see
Section 23.3. This includes string representations of XML values, such as in the above examples. This
would ordinarily mean that encoding declarations contained in XML data can become invalid as the
character data is converted to other encodings while traveling between client and server, because the
embedded encoding declaration is not changed. To cope with this behavior, encoding declarations
contained in character strings presented for input to the xml type are ignored, and content is assumed
to be in the current server encoding. Consequently, for correct processing, character strings of XML
data must be sent from the client in the current client encoding. It is the responsibility of the client
to either convert documents to the current client encoding before sending them to the server, or to
adjust the client encoding appropriately. On output, values of type xml will not have an encoding
declaration, and clients should assume all data is in the current client encoding.

When using binary mode to pass query parameters to the server and query results back to the client, no
encoding conversion is performed, so the situation is different. In this case, an encoding declaration
in the XML data will be observed, and if it is absent, the data will be assumed to be in UTF-8 (as
required by the XML standard; note that PostgreSQL does not support UTF-16). On output, data will
Data Types

have an encoding declaration specifying the client encoding, unless the client encoding is UTF-8, in which case it will be omitted.

 Needless to say, processing XML data with PostgreSQL will be less error-prone and more efficient if the XML data encoding, client encoding, and server encoding are the same. Since XML data is internally processed in UTF-8, computations will be most efficient if the server encoding is also UTF-8.

Caution

Some XML-related functions may not work at all on non-ASCII data when the server encoding is not UTF-8. This is known to be an issue for xmltable() and xpath() in particular.

8.13.3. Accessing XML Values

The xml data type is unusual in that it does not provide any comparison operators. This is because there is no well-defined and universally useful comparison algorithm for XML data. One consequence of this is that you cannot retrieve rows by comparing an xml column against a search value. XML values should therefore typically be accompanied by a separate key field such as an ID. An alternative solution for comparing XML values is to convert them to character strings first, but note that character string comparison has little to do with a useful XML comparison method.

Since there are no comparison operators for the xml data type, it is not possible to create an index directly on a column of this type. If speedy searches in XML data are desired, possible workarounds include casting the expression to a character string type and indexing that, or indexing an XPath expression. Of course, the actual query would have to be adjusted to search by the indexed expression.

The text-search functionality in PostgreSQL can also be used to speed up full-document searches of XML data. The necessary preprocessing support is, however, not yet available in the PostgreSQL distribution.

8.14. JSON Types

JSON data types are for storing JSON (JavaScript Object Notation) data, as specified in RFC 7159\(^1\). Such data can also be stored as text, but the JSON data types have the advantage of enforcing that each stored value is valid according to the JSON rules. There are also assorted JSON-specific functions and operators available for data stored in these data types; see Section 9.15.

There are two JSON data types: json and jsonb. They accept almost identical sets of values as input. The major practical difference is one of efficiency. The json data type stores an exact copy of the input text, which processing functions must reparse on each execution; while jsonb data is stored in a decomposed binary format that makes it slightly slower to input due to added conversion overhead, but significantly faster to process, since no reparsing is needed. jsonb also supports indexing, which can be a significant advantage.

Because the json type stores an exact copy of the input text, it will preserve semantically-insignificant white space between tokens, as well as the order of keys within JSON objects. Also, if a JSON object within the value contains the same key more than once, all the key/value pairs are kept. (The processing functions consider the last value as the operative one.) By contrast, jsonb does not preserve white space, does not preserve the order of object keys, and does not keep duplicate object keys. If duplicate keys are specified in the input, only the last value is kept.

In general, most applications should prefer to store JSON data as jsonb, unless there are quite specialized needs, such as legacy assumptions about ordering of object keys.

---

\(^1\) https://tools.ietf.org/html/rfc7159
PostgreSQL allows only one character set encoding per database. It is therefore not possible for the JSON types to conform rigidly to the JSON specification unless the database encoding is UTF8. Attempts to directly include characters that cannot be represented in the database encoding will fail; conversely, characters that can be represented in the database encoding but not in UTF8 will be allowed.

RFC 7159 permits JSON strings to contain Unicode escape sequences denoted by \uXXXX. In the input function for the json type, Unicode escapes are allowed regardless of the database encoding, and are checked only for syntactic correctness (that is, that four hex digits follow \u). However, the input function for jsonb is stricter: it disallows Unicode escapes for non-ASCII characters (those above U+007F) unless the database encoding is UTF8. The jsonb type also rejects \u0000 (because that cannot be represented in PostgreSQL's text type), and it insists that any use of Unicode surrogate pairs to designate characters outside the Unicode Basic Multilingual Plane be correct. Valid Unicode escapes are converted to the equivalent ASCII or UTF8 character for storage; this includes folding surrogate pairs into a single character.

**Note**

Many of the JSON processing functions described in Section 9.15 will convert Unicode escapes to regular characters, and will therefore throw the same types of errors just described even if their input is of type json not jsonb. The fact that the json input function does not make these checks may be considered a historical artifact, although it does allow for simple storage (without processing) of JSON Unicode escapes in a non-UTF8 database encoding. In general, it is best to avoid mixing Unicode escapes in JSON with a non-UTF8 database encoding, if possible.

When converting textual JSON input into jsonb, the primitive types described by RFC 7159 are effectively mapped onto native PostgreSQL types, as shown in Table 8.23. Therefore, there are some minor additional constraints on what constitutes valid jsonb data that do not apply to the json type, nor to JSON in the abstract, corresponding to limits on what can be represented by the underlying data type. Notably, jsonb will reject numbers that are outside the range of the PostgreSQL numeric data type, while json will not. Such implementation-defined restrictions are permitted by RFC 7159. However, in practice such problems are far more likely to occur in other implementations, as it is common to represent JSON’s number primitive type as IEEE 754 double precision floating point (which RFC 7159 explicitly anticipates and allows for). When using JSON as an interchange format with such systems, the danger of losing numeric precision compared to data originally stored by PostgreSQL should be considered.

Conversely, as noted in the table there are some minor restrictions on the input format of JSON primitive types that do not apply to the corresponding PostgreSQL types.

<table>
<thead>
<tr>
<th>JSON primitive type</th>
<th>PostgreSQL type</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>string</td>
<td>text</td>
<td>\u0000 is disallowed, as are non-ASCII Unicode escapes if database encoding is not UTF8</td>
</tr>
<tr>
<td>number</td>
<td>numeric</td>
<td>NaN and infinity values are disallowed</td>
</tr>
<tr>
<td>boolean</td>
<td>boolean</td>
<td>Only lowercase true and false spellings are accepted</td>
</tr>
<tr>
<td>null</td>
<td>(none)</td>
<td>SQL NULL is a different concept</td>
</tr>
</tbody>
</table>

## 8.14.1. JSON Input and Output Syntax

The input/output syntax for the JSON data types is as specified in RFC 7159.
Data Types

The following are all valid json (or jsonb) expressions:

-- Simple scalar/primitive value
-- Primitive values can be numbers, quoted strings, true, false, or null
SELECT '5'::json;

-- Array of zero or more elements (elements need not be of same type)
SELECT '[1, 2, "foo", null]'::json;

-- Object containing pairs of keys and values
-- Note that object keys must always be quoted strings
SELECT '{"bar": "baz", "balance": 7.77, "active": false}'::json;

-- Arrays and objects can be nested arbitrarily
SELECT '{"foo": [{true, "bar"}, {"tags":{"a": 1, "b": null}}]}'::json;

As previously stated, when a JSON value is input and then printed without any additional processing, json outputs the same text that was input, while jsonb does not preserve semantically-insignificant details such as whitespace. For example, note the differences here:

SELECT '{"bar": "baz", "balance": 7.77, "active":false}'::json;
{
  "bar": "baz",
  "balance": 7.77,
  "active":false
}(1 row)

SELECT '{"bar": "baz", "balance": 7.77, "active":false}'::jsonb;
{
  "bar": "baz",
  "active": false,
  "balance": 7.77
}(1 row)

One semantically-insignificant detail worth noting is that in jsonb, numbers will be printed according to the behavior of the underlying numeric type. In practice this means that numbers entered with E notation will be printed without it, for example:

SELECT '{"reading": 1.230e-5}':'json, '{"reading": 1.230e-5}':'jsonb;
{
  "reading": 1.230e-5
} | {
  "reading": 0.00001230
}(1 row)

However, jsonb will preserve trailing fractional zeroes, as seen in this example, even though those are semantically insignificant for purposes such as equality checks.

8.14.2. Designing JSON documents effectively

Representing data as JSON can be considerably more flexible than the traditional relational data model, which is compelling in environments where requirements are fluid. It is quite possible for both approaches to co-exist and complement each other within the same application. However, even for applications where maximal flexibility is desired, it is still recommended that JSON documents have a somewhat fixed structure. The structure is typically unenforced (though enforcing some business
Data Types

rules declaratively is possible), but having a predictable structure makes it easier to write queries that usefully summarize a set of “documents” (datums) in a table.

JSON data is subject to the same concurrency-control considerations as any other data type when stored in a table. Although storing large documents is practicable, keep in mind that any update acquires a row-level lock on the whole row. Consider limiting JSON documents to a manageable size in order to decrease lock contention among updating transactions. Ideally, JSON documents should each represent an atomic datum that business rules dictate cannot reasonably be further subdivided into smaller datums that could be modified independently.

8.14.3. jsonb Containment and Existence

Testing containment is an important capability of jsonb. There is no parallel set of facilities for the json type. Containment tests whether one jsonb document has contained within it another one. These examples return true except as noted:

-- Simple scalar/primitive values contain only the identical value:
SELECT "foo"::jsonb @> "foo"::jsonb;

-- The array on the right side is contained within the one on the left:
SELECT '[1, 2, 3]':jsonb @> '[1, 3]':jsonb;

-- Order of array elements is not significant, so this is also true:
SELECT '[1, 2, 3]':jsonb @> '[3, 1]':jsonb;

-- Duplicate array elements don't matter either:
SELECT '[1, 2, 3]':jsonb @> '[1, 2, 2]':jsonb;

-- The object with a single pair on the right side is contained
-- within the object on the left side:
SELECT '{"product": "PostgreSQL", "version": 9.4, "jsonb":
true}':jsonb @> '{"version": 9.4}':jsonb;

-- The array on the right side is not considered contained within
-- array on the left, even though a similar array is nested within it:
SELECT '[1, 2, [1, 3]]':jsonb @> '[1, 3]':jsonb; -- yields false

-- But with a layer of nesting, it is contained:
SELECT '[1, 2, [1, 3]]':jsonb @> '[[1, 3]]':jsonb;

-- Similarly, containment is not reported here:
SELECT '{"foo": {"bar": "baz"}}':jsonb @> '{"bar": "baz"}':jsonb;

-- A top-level key and an empty object is contained:
SELECT '{"foo": {"bar": "baz"}}':jsonb @> '{"foo": {}}':jsonb;

The general principle is that the contained object must match the containing object as to structure and data contents, possibly after discarding some non-matching array elements or object key/value pairs from the containing object. But remember that the order of array elements is not significant when doing a containment match, and duplicate array elements are effectively considered only once.

As a special exception to the general principle that the structures must match, an array may contain a primitive value:
Data Types

-- This array contains the primitive string value:
SELECT '["foo", "bar"]'::jsonb @> '"bar"'::jsonb;
-- This exception is not reciprocal -- non-containment is reported
here:
SELECT '"bar"'::jsonb @> '["bar"]'::jsonb; -- yields false
jsonb also has an existence operator, which is a variation on the theme of containment: it tests
whether a string (given as a text value) appears as an object key or array element at the top level of
the jsonb value. These examples return true except as noted:
-- String exists as array element:
SELECT '["foo", "bar", "baz"]'::jsonb ? 'bar';
-- String exists as object key:
SELECT '{"foo": "bar"}'::jsonb ? 'foo';
-- Object values are not considered:
SELECT '{"foo": "bar"}'::jsonb ? 'bar';

-- yields false

-- As with containment, existence must match at the top level:
SELECT '{"foo": {"bar": "baz"}}'::jsonb ? 'bar'; -- yields false
-- A string is considered to exist if it matches a primitive JSON
string:
SELECT '"foo"'::jsonb ? 'foo';
JSON objects are better suited than arrays for testing containment or existence when there are many
keys or elements involved, because unlike arrays they are internally optimized for searching, and do
not need to be searched linearly.

Tip
Because JSON containment is nested, an appropriate query can skip explicit selection
of sub-objects. As an example, suppose that we have a doc column containing objects at the top level, with most objects containing tags fields that contain arrays
of sub-objects. This query finds entries in which sub-objects containing both "term":"paris" and "term":"food" appear, while ignoring any such keys outside
the tags array:
SELECT doc->'site_name' FROM websites
WHERE doc @> '{"tags":[{"term":"paris"},
{"term":"food"}]}';
One could accomplish the same thing with, say,
SELECT doc->'site_name' FROM websites
WHERE doc->'tags' @> '[{"term":"paris"},
{"term":"food"}]';
but that approach is less flexible, and often less efficient as well.
On the other hand, the JSON existence operator is not nested: it will only look for the
specified key or array element at top level of the JSON value.

169


The various containment and existence operators, along with all other JSON operators and functions are documented in Section 9.15.

### 8.14.4. `jsonb` Indexing

GIN indexes can be used to efficiently search for keys or key/value pairs occurring within a large number of `jsonb` documents (datums). Two GIN “operator classes” are provided, offering different performance and flexibility trade-offs.

The default GIN operator class for `jsonb` supports queries with top-level key-exists operators `?`, `?&` and `? |` operators and path/value-exists operator `@>`. (For details of the semantics that these operators implement, see Table 9.44.) An example of creating an index with this operator class is:

```sql
CREATE INDEX idxgin ON api USING GIN (jdoc);
```

The non-default GIN operator class `jsonb_path_ops` supports indexing the `@>` operator only. An example of creating an index with this operator class is:

```sql
CREATE INDEX idxginp ON api USING GIN (jdoc jsonb_path_ops);
```

Consider the example of a table that stores JSON documents retrieved from a third-party web service, with a documented schema definition. A typical document is:

```
{
  "guid": "9c36adcl-7fb5-4d5b-83b4-90356a46061a",
  "name": "Angela Barton",
  "is_active": true,
  "company": "Magnafone",
  "address": "178 Howard Place, Gulf, Washington, 702",
  "registered": "2009-11-07T08:53:22 +08:00",
  "latitude": 19.793713,
  "longitude": 86.513373,
  "tags": ["enim", "aliquip", "qui"
    ]
}
```

We store these documents in a table named `api`, in a `jsonb` column named `jdoc`. If a GIN index is created on this column, queries like the following can make use of the index:

```sql
-- Find documents in which the key "company" has value "Magnafone"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc @> '{"company": "Magnafone"}';
```

However, the index could not be used for queries like the following, because though the operator `?` is indexable, it is not applied directly to the indexed column `jdoc`:

```sql
-- Find documents in which the key "tags" contains key or array element "qui"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc -> 'tags' ? 'qui';
```

Still, with appropriate use of expression indexes, the above query can use an index. If querying for particular items within the "tags" key is common, defining an index like this may be worthwhile:
CREATE INDEX idxgintags ON api USING GIN ((jdoc -> 'tags'));

Now, the WHERE clause jdoc -> 'tags' ? 'qui' will be recognized as an application of the indexable operator ? to the indexed expression jdoc -> 'tags'. (More information on expression indexes can be found in Section 11.7.)

Another approach to querying is to exploit containment, for example:

-- Find documents in which the key "tags" contains array element "qui"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc @> '{"tags": ["qui"]}';

A simple GIN index on the jdoc column can support this query. But note that such an index will store copies of every key and value in the jdoc column, whereas the expression index of the previous example stores only data found under the tags key. While the simple-index approach is far more flexible (since it supports queries about any key), targeted expression indexes are likely to be smaller and faster to search than a simple index.

Although the jsonb_path_ops operator class supports only queries with the @> operator, it has notable performance advantages over the default operator class jsonb_ops. A jsonb_path_ops index is usually much smaller than a jsonb_ops index over the same data, and the specificity of searches is better, particularly when queries contain keys that appear frequently in the data. Therefore search operations typically perform better than with the default operator class.

The technical difference between a jsonb_ops and a jsonb_path_ops GIN index is that the former creates independent index items for each key and value in the data, while the latter creates index items only for each value in the data. ² Basically, each jsonb_path_ops index item is a hash of the value and the key(s) leading to it; for example to index {"foo": {"bar": "baz"}}, a single index item would be created incorporating all three of foo, bar, and baz into the hash value. Thus a containment query looking for this structure would result in an extremely specific index search; but there is no way at all to find out whether foo appears as a key. On the other hand, a jsonb_ops index would create three index items representing foo, bar, and baz separately; then to do the containment query, it would look for rows containing all three of these items. While GIN indexes can perform such an AND search fairly efficiently, it will still be less specific and slower than the equivalent jsonb_path_ops search, especially if there are a very large number of rows containing any single one of the three index items.

A disadvantage of the jsonb_path_ops approach is that it produces no index entries for JSON structures not containing any values, such as {"a": {}}. If a search for documents containing such a structure is requested, it will require a full-index scan, which is quite slow. jsonb_path_ops is therefore ill-suited for applications that often perform such searches.

jsonb also supports btree and hash indexes. These are usually useful only if it's important to check equality of complete JSON documents. The btree ordering for jsonb datums is seldom of great interest, but for completeness it is:

Object > Array > Boolean > Number > String > Null
Object with n pairs > object with n - 1 pairs
Array with n elements > array with n - 1 elements
Objects with equal numbers of pairs are compared in the order:

² For this purpose, the term “value” includes array elements, though JSON terminology sometimes considers array elements distinct from values within objects.
key-1, value-1, key-2 ...

Note that object keys are compared in their storage order; in particular, since shorter keys are stored before longer keys, this can lead to results that might be unintuitive, such as:

{ "aa": 1, "c": 1} > {"b": 1, "d": 1}

Similarly, arrays with equal numbers of elements are compared in the order:

element-1, element-2 ...

Primitive JSON values are compared using the same comparison rules as for the underlying PostgreSQL data type. Strings are compared using the default database collation.

8.14.5. Transforms

Additional extensions are available that implement transforms for the jsonb type for different procedural languages.

The extensions for PL/Perl are called jsonb_plperl and jsonb_plperlu. If you use them, jsonb values are mapped to Perl arrays, hashes, and scalars, as appropriate.

The extensions for PL/Python are called jsonb_plpythonu, jsonb_plpython2u, and jsonb_plpython3u (see Section 46.1 for the PL/Python naming convention). If you use them, jsonb values are mapped to Python dictionaries, lists, and scalars, as appropriate.

8.15. Arrays

PostgreSQL allows columns of a table to be defined as variable-length multidimensional arrays. Arrays of any built-in or user-defined base type, enum type, composite type, range type, or domain can be created.

8.15.1. Declaration of Array Types

To illustrate the use of array types, we create this table:

```
CREATE TABLE sal_emp ( 
    name text, 
    pay_by_quarter integer[], 
    schedule text[][] 
); 
```

As shown, an array data type is named by appending square brackets ([]) to the data type name of the array elements. The above command will create a table named sal_emp with a column of type text (name), a one-dimensional array of type integer (pay_by_quarter), which represents the employee's salary by quarter, and a two-dimensional array of text (schedule), which represents the employee's weekly schedule.

The syntax for CREATE TABLE allows the exact size of arrays to be specified, for example:

```
CREATE TABLE tictactoe ( 
    squares integer[3][3] 
) 
```
However, the current implementation ignores any supplied array size limits, i.e., the behavior is the same as for arrays of unspecified length.

The current implementation does not enforce the declared number of dimensions either. Arrays of a particular element type are all considered to be of the same type, regardless of size or number of dimensions. So, declaring the array size or number of dimensions in CREATE TABLE is simply documentation; it does not affect run-time behavior.

An alternative syntax, which conforms to the SQL standard by using the keyword ARRAY, can be used for one-dimensional arrays. pay_by_quarter could have been defined as:

```sql
pay_by_quarter  integer ARRAY[4],
```

Or, if no array size is to be specified:

```sql
pay_by_quarter  integer ARRAY,
```

As before, however, PostgreSQL does not enforce the size restriction in any case.

### 8.15.2. Array Value Input

To write an array value as a literal constant, enclose the element values within curly braces and separate them by commas. (If you know C, this is not unlike the C syntax for initializing structures.) You can put double quotes around any element value, and must do so if it contains commas or curly braces. (More details appear below.) Thus, the general format of an array constant is the following:

```
'{ val1 delim val2 delim ... }'
```

where `delim` is the delimiter character for the type, as recorded in its `pg_type` entry. Among the standard data types provided in the PostgreSQL distribution, all use a comma (,), except for type `box` which uses a semicolon (;). Each `val` is either a constant of the array element type, or a subarray. An example of an array constant is:

```
'{ {1,2,3}, {4,5,6}, {7,8,9} }'
```

This constant is a two-dimensional, 3-by-3 array consisting of three subarrays of integers.

To set an element of an array constant to NULL, write `NULL` for the element value. (Any upper- or lower-case variant of `NULL` will do.) If you want an actual string value “NULL”, you must put double quotes around it.

(These kinds of array constants are actually only a special case of the generic type constants discussed in Section 4.1.2.7. The constant is initially treated as a string and passed to the array input conversion routine. An explicit type specification might be necessary.)

Now we can show some INSERT statements:

```sql
INSERT INTO sal_emp
VALUES ('Bill',
'{10000, 10000, 10000, 10000}',
'{{"meeting", "lunch"}, {"training", "presentation"}}');
```
INSERT INTO sal_emp
VALUES ('Carol',
    '{20000, 25000, 25000, 25000}',
    '{"breakfast", "consulting"},
    {"meeting", "lunch"});

The result of the previous two inserts looks like this:

```
SELECT * FROM sal_emp;
name |      pay_by_quarter       |                 schedule
-------+---------------------------|-------------------------
  Bill | {10000,10000,10000,10000} | {{meeting,lunch},
    {training,presentation}}
 Carol | {20000,25000,25000,25000} | {{breakfast,consulting},
    {meeting,lunch}}
```

Multidimensional arrays must have matching extents for each dimension. A mismatch causes an error, for example:

```
INSERT INTO sal_emp
VALUES ('Bill',
    '{10000, 10000, 10000, 10000}',
    {{"meeting", "lunch"}, {"meeting"}});
ERROR:  multidimensional arrays must have array expressions with
        matching dimensions
```

The `ARRAY` constructor syntax can also be used:

```
INSERT INTO sal_emp
VALUES ('Bill',
    ARRAY[10000, 10000, 10000, 10000],
    ARRAY[['meeting', 'lunch'], ['training', 'presentation']]);
```

```
INSERT INTO sal_emp
VALUES ('Carol',
    ARRAY[20000, 25000, 25000, 25000],
    ARRAY[['breakfast', 'consulting'], ['meeting', 'lunch']]));
```

Notice that the array elements are ordinary SQL constants or expressions; for instance, string literals are single quoted, instead of double quoted as they would be in an array literal. The `ARRAY` constructor syntax is discussed in more detail in Section 4.2.12.

### 8.15.3. Accessing Arrays

Now, we can run some queries on the table. First, we show how to access a single element of an array. This query retrieves the names of the employees whose pay changed in the second quarter:

```
SELECT name FROM sal_emp WHERE pay_by_quarter[1] <>
    pay_by_quarter[2];
```

```
 name
-------
 Carol
(1 row)
```
Data Types

The array subscript numbers are written within square brackets. By default PostgreSQL uses a one-based numbering convention for arrays, that is, an array of \( n \) elements starts with \( \text{array}[1] \) and ends with \( \text{array}[n] \).

This query retrieves the third quarter pay of all employees:

```
SELECT pay_by_quarter[3] FROM sal_emp;
```

<table>
<thead>
<tr>
<th>pay_by_quarter</th>
</tr>
</thead>
<tbody>
<tr>
<td>10000</td>
</tr>
<tr>
<td>25000</td>
</tr>
</tbody>
</table>

(2 rows)

We can also access arbitrary rectangular slices of an array, or subarrays. An array slice is denoted by writing \( \text{lower-bound:upper-bound} \) for one or more array dimensions. For example, this query retrieves the first item on Bill’s schedule for the first two days of the week:

```
SELECT schedule[1:2][1:1] FROM sal_emp WHERE name = 'Bill';
```

<table>
<thead>
<tr>
<th>schedule</th>
</tr>
</thead>
<tbody>
<tr>
<td>{{meeting},{training}}</td>
</tr>
</tbody>
</table>

(1 row)

If any dimension is written as a slice, i.e., contains a colon, then all dimensions are treated as slices. Any dimension that has only a single number (no colon) is treated as being from 1 to the number specified. For example, \( [2] \) is treated as \( [1:2] \), as in this example:

```
SELECT schedule[1:2][2] FROM sal_emp WHERE name = 'Bill';
```

<table>
<thead>
<tr>
<th>schedule</th>
</tr>
</thead>
<tbody>
<tr>
<td>{{meeting,lunch},{training,presentation}}</td>
</tr>
</tbody>
</table>

(1 row)

To avoid confusion with the non-slice case, it’s best to use slice syntax for all dimensions, e.g., \( [1:2] [1:1] \), not \( [2][1:1] \).

It is possible to omit the \( \text{lower-bound} \) and/or \( \text{upper-bound} \) of a slice specifier; the missing bound is replaced by the lower or upper limit of the array’s subscripts. For example:

```
SELECT schedule[:2][2:] FROM sal_emp WHERE name = 'Bill';
```

<table>
<thead>
<tr>
<th>schedule</th>
</tr>
</thead>
<tbody>
<tr>
<td>{{lunch},{presentation}}</td>
</tr>
</tbody>
</table>

(1 row)

```
SELECT schedule[:][1:1] FROM sal_emp WHERE name = 'Bill';
```

<table>
<thead>
<tr>
<th>schedule</th>
</tr>
</thead>
<tbody>
<tr>
<td>{{meeting},{training}}</td>
</tr>
</tbody>
</table>

(1 row)
An array subscript expression will return null if either the array itself or any of the subscript expressions are null. Also, null is returned if a subscript is outside the array bounds (this case does not raise an error). For example, if \textit{schedule} currently has the dimensions \([1:3][1:2]\) then referencing \texttt{schedule[3][3]} yields NULL. Similarly, an array reference with the wrong number of subscripts yields a null rather than an error.

An array slice expression likewise yields null if the array itself or any of the subscript expressions are null. However, in other cases such as selecting an array slice that is completely outside the current array bounds, a slice expression yields an empty (zero-dimensional) array instead of null. (This does not match non-slice behavior and is done for historical reasons.) If the requested slice partially overlaps the array bounds, then it is silently reduced to just the overlapping region instead of returning null.

The current dimensions of any array value can be retrieved with the \texttt{array_dims} function:

\begin{verbatim}
SELECT array_dims(schedule) FROM sal_emp WHERE name = 'Carol';
array_dims
------------
[1:2][1:2]
(1 row)
\end{verbatim}

\texttt{array_dims} produces a text result, which is convenient for people to read but perhaps inconvenient for programs. Dimensions can also be retrieved with \texttt{array_upper} and \texttt{array_lower}, which return the upper and lower bound of a specified array dimension, respectively:

\begin{verbatim}
SELECT array_upper(schedule, 1) FROM sal_emp WHERE name = 'Carol';
array_upper
----------
2
(1 row)
\end{verbatim}

\texttt{array_length} will return the length of a specified array dimension:

\begin{verbatim}
SELECT array_length(schedule, 1) FROM sal_emp WHERE name = 'Carol';
array_length
----------
2
(1 row)
\end{verbatim}

\texttt{cardinality} returns the total number of elements in an array across all dimensions. It is effectively the number of rows a call to \texttt{unnest} would yield:

\begin{verbatim}
SELECT cardinality(schedule) FROM sal_emp WHERE name = 'Carol';
cardinality
----------
4
(1 row)
\end{verbatim}

\textbf{8.15.4. Modifying Arrays}

An array value can be replaced completely:
UPDATE sal_emp SET pay_by_quarter = '{25000,25000,27000,27000}'
  WHERE name = 'Carol';

or using the ARRAY expression syntax:

UPDATE sal_emp SET pay_by_quarter = ARRAY[25000,25000,27000,27000]
  WHERE name = 'Carol';

An array can also be updated at a single element:

  WHERE name = 'Bill';

or updated in a slice:

UPDATE sal_emp SET pay_by_quarter[1:2] = '{27000,27000}'
  WHERE name = 'Carol';

The slice syntaxes with omitted lower-bound and/or upper-bound can be used too, but only when updating an array value that is not NULL or zero-dimensional (otherwise, there is no existing subscript limit to substitute).

A stored array value can be enlarged by assigning to elements not already present. Any positions between those previously present and the newly assigned elements will be filled with nulls. For example, if array myarray currently has 4 elements, it will have six elements after an update that assigns to myarray[6]:myarray[5] will contain null. Currently, enlargement in this fashion is only allowed for one-dimensional arrays, not multidimensional arrays.

Subscripted assignment allows creation of arrays that do not use one-based subscripts. For example one might assign to myarray[-2:7] to create an array with subscript values from -2 to 7.

New array values can also be constructed using the concatenation operator, ||:

SELECT ARRAY[1,2] || ARRAY[3,4];
?column?
-----------
{1,2,3,4}
(1 row)

SELECT ARRAY[5,6] || ARRAY[[1,2],[3,4]];
?column?
------------------
{{5,6},{1,2},{3,4}}
(1 row)

The concatenation operator allows a single element to be pushed onto the beginning or end of a one-dimensional array. It also accepts two $N$-dimensional arrays, or an $N$-dimensional and an $N+1$-dimensional array.

When a single element is pushed onto either the beginning or end of a one-dimensional array, the result is an array with the same lower bound subscript as the array operand. For example:

SELECT array_dims(1 || '0:1={2,3}':::int[]);
array_dims
-----------
When two arrays with an equal number of dimensions are concatenated, the result retains the lower bound subscript of the left-hand operand's outer dimension. The result is an array comprising every element of the left-hand operand followed by every element of the right-hand operand. For example:

```
SELECT array_dims(ARRAY[1, 2] || ARRAY[3, 4, 5]);
array_dims
------------
[1:5]
(1 row)
```

When an \(N\)-dimensional array is pushed onto the beginning or end of an \(N+1\)-dimensional array, the result is analogous to the element-array case above. Each \(N\)-dimensional sub-array is essentially an element of the \(N+1\)-dimensional array's outer dimension. For example:

```
SELECT array_dims(ARRAY[1, 2] || ARRAY[[3, 4], [5, 6]]);
array_dims
------------
[1:3][1:2]
(1 row)
```

An array can also be constructed by using the functions `array_prepend`, `array_append`, or `array_cat`. The first two only support one-dimensional arrays, but `array_cat` supports multidimensional arrays. Some examples:

```
SELECT array_prepend(1, ARRAY[2, 3]);
array_prepend
-------------
{1, 2, 3}
(1 row)
```

```
SELECT array_append(ARRAY[1, 2], 3);
array_append
-------------
{1, 2, 3}
(1 row)
```

```
SELECT array_cat(ARRAY[1, 2], ARRAY[3, 4]);
array_cat
----------
{1, 2, 3, 4}
(1 row)
```
SELECT array_cat(ARRAY[[1,2],[3,4]], ARRAY[5,6]);
array_cat
---------------------
{{1,2},{3,4},{5,6}}
(1 row)

SELECT array_cat(ARRAY[5,6], ARRAY[[1,2],[3,4]]);
array_cat
---------------------
{{5,6},{1,2},{3,4}}

In simple cases, the concatenation operator discussed above is preferred over direct use of these functions. However, because the concatenation operator is overloaded to serve all three cases, there are situations where use of one of the functions is helpful to avoid ambiguity. For example consider:

SELECT ARRAY[1, 2] || '{3, 4}';  -- the untyped literal is taken as an array
?column?
-----------
{1,2,3,4}

SELECT ARRAY[1, 2] || '7';                 -- so is this one
ERROR:  malformed array literal: "7"

SELECT ARRAY[1, 2] || NULL;                -- so is an undecorated NULL
?column?
----------
{1,2}
(1 row)

SELECT array_append(ARRAY[1, 2], NULL);    -- this might have been meant
array_append
-------------
{1,2,NULL}

In the examples above, the parser sees an integer array on one side of the concatenation operator, and a constant of undetermined type on the other. The heuristic it uses to resolve the constant's type is to assume it's of the same type as the operator's other input — in this case, integer array. So the concatenation operator is presumed to represent array_cat, not array_append. When that's the wrong choice, it could be fixed by casting the constant to the array's element type; but explicit use of array_append might be a preferable solution.

8.15.5. Searching in Arrays

To search for a value in an array, each value must be checked. This can be done manually, if you know the size of the array. For example:


However, this quickly becomes tedious for large arrays, and is not helpful if the size of the array is unknown. An alternative method is described in Section 9.23. The above query could be replaced by:
SELECT * FROM sal_emp WHERE 10000 = ANY (pay_by_quarter);
In addition, you can find rows where the array has all values equal to 10000 with:

SELECT * FROM sal_emp WHERE 10000 = ALL (pay_by_quarter);
Alternatively, the `generate_subscripts` function can be used. For example:

SELECT * FROM
  (SELECT pay_by_quarter,
   generate_subscripts(pay_by_quarter, 1) AS s
  FROM sal_emp) AS foo
WHERE pay_by_quarter[s] = 10000;
This function is described in Table 9.59.

You can also search an array using the `&&` operator, which checks whether the left operand overlaps
with the right operand. For instance:

SELECT * FROM sal_emp WHERE pay_by_quarter && ARRAY[10000];
This and other array operators are further described in Section 9.18. It can be accelerated by an
appropriate index, as described in Section 11.2.

You can also search for specific values in an array using the `array_position` and `array_positions`
functions. The former returns the subscript of the first occurrence of a value in an array;
the latter returns an array with the subscripts of all occurrences of the value in the array. For example:

SELECT
  array_position(ARRAY['sun','mon','tue','wed','thu','fri','sat'],
  'mon');
array_positions
------------------
2
SELECT array_positions(ARRAY[1, 4, 3, 1, 3, 4, 2, 1], 1);
array_positions
------------------
{1,4,8}

Tip

Arrays are not sets; searching for specific array elements can be a sign of database
misdesign. Consider using a separate table with a row for each item that would be an
array element. This will be easier to search, and is likely to scale better for a large
number of elements.

8.15.6. Array Input and Output Syntax

The external text representation of an array value consists of items that are interpreted according to
the I/O conversion rules for the array's element type, plus decoration that indicates the array structure.
The decoration consists of curly braces ({ and }) around the array value plus delimiter characters
between adjacent items. The delimiter character is usually a comma (,) but can be something else: it
Data Types

is determined by the typdelim setting for the array’s element type. Among the standard data types provided in the PostgreSQL distribution, all use a comma, except for type box, which uses a semicolon (;). In a multidimensional array, each dimension (row, plane, cube, etc.) gets its own level of curly braces, and delimiters must be written between adjacent curly-braced entities of the same level.

The array output routine will put double quotes around element values if they are empty strings, contain curly braces, delimiter characters, double quotes, backslashes, or white space, or match the word NULL. Double quotes and backslashes embedded in element values will be backslash-escaped. For numeric data types it is safe to assume that double quotes will never appear, but for textual data types one should be prepared to cope with either the presence or absence of quotes.

By default, the lower bound index value of an array’s dimensions is set to one. To represent arrays with other lower bounds, the array subscript ranges can be specified explicitly before writing the array contents. This decoration consists of square brackets ([ ]) around each array dimension’s lower and upper bounds, with a colon (:) delimiter character in between. The array dimension decoration is followed by an equal sign (=). For example:

```sql
SELECT f1[1][2][3] AS e1, f1[1][-1][5] AS e2
FROM (SELECT '{[1:1][-2:-1][3:5]={{1,2,3},{4,5,6}}}'::int[] AS f1)
AS ss;
```

<table>
<thead>
<tr>
<th>e1</th>
<th>e2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>6</td>
</tr>
</tbody>
</table>

The array output routine will include explicit dimensions in its result only when there are one or more lower bounds different from one.

If the value written for an element is NULL (in any case variant), the element is taken to be NULL. The presence of any quotes or backslashes disables this and allows the literal string value “NULL” to be entered. Also, for backward compatibility with pre-8.2 versions of PostgreSQL, the array_nulls configuration parameter can be turned off to suppress recognition of NULL as a NULL.

As shown previously, when writing an array value you can use double quotes around any individual array element. You must do so if the element value would otherwise confuse the array-value parser. For example, elements containing curly braces, commas (or the data type’s delimiter character), double quotes, backslashes, or leading or trailing whitespace must be double-quoted. Empty strings and strings matching the word NULL must be quoted, too. To put a double quote or backslash in a quoted array element value, precede it with a backslash. Alternatively, you can avoid quotes and use backslash-escaping to protect all data characters that would otherwise be taken as array syntax.

You can add whitespace before a left brace or after a right brace. You can also add whitespace before or after any individual item string. In all of these cases the whitespace will be ignored. However, whitespace within double-quoted elements, or surrounded on both sides by non-whitespace characters of an element, is not ignored.

Tip

The ARRAY constructor syntax (see Section 4.2.12) is often easier to work with than the array-literal syntax when writing array values in SQL commands. In ARRAY, individual element values are written the same way they would be written when not members of an array.

8.16. Composite Types
A composite type represents the structure of a row or record; it is essentially just a list of field names and their data types. PostgreSQL allows composite types to be used in many of the same ways that simple types can be used. For example, a column of a table can be declared to be of a composite type.

### 8.16.1. Declaration of Composite Types

Here are two simple examples of defining composite types:

```sql
CREATE TYPE complex AS (  
r       double precision,  
i       double precision
);
CREATE TYPE inventory_item AS (  
name            text,  
supplier_id     integer,  
price           numeric
);
```

The syntax is comparable to `CREATE TABLE`, except that only field names and types can be specified; no constraints (such as `NOT NULL`) can presently be included. Note that the `AS` keyword is essential; without it, the system will think a different kind of `CREATE TYPE` command is meant, and you will get odd syntax errors.

Having defined the types, we can use them to create tables:

```sql
CREATE TABLE on_hand (  
    item      inventory_item,  
    count     integer
);
```

```sql
INSERT INTO on_hand VALUES (ROW('fuzzy dice', 42, 1.99), 1000);
```

or functions:

```sql
CREATE FUNCTION price_extension(inventory_item, integer) RETURNS numeric  
AS 'SELECT $1.price * $2' LANGUAGE SQL;
```

```sql
SELECT price_extension(item, 10) FROM on_hand;
```

Whenever you create a table, a composite type is also automatically created, with the same name as the table, to represent the table's row type. For example, had we said:

```sql
CREATE TABLE inventory_item (  
    name            text,  
    supplier_id     integer REFERENCES suppliers,  
    price           numeric CHECK (price > 0)
);
```

then the same `inventory_item` composite type shown above would come into being as a byproduct, and could be used just as above. Note however an important restriction of the current implementation: since no constraints are associated with a composite type, the constraints shown in the table definition do not apply to values of the composite type outside the table. (To work around this, create a domain over the composite type, and apply the desired constraints as `CHECK` constraints of the domain.)
8.16.2. Constructing Composite Values

To write a composite value as a literal constant, enclose the field values within parentheses and separate them by commas. You can put double quotes around any field value, and must do so if it contains commas or parentheses. (More details appear below.) Thus, the general format of a composite constant is the following:

' ( val1 , val2 , ... ) '

An example is:

' ("fuzzy dice",42,1.99) '

which would be a valid value of the inventory_item type defined above. To make a field be NULL, write no characters at all in its position in the list. For example, this constant specifies a NULL third field:

' ("fuzzy dice",42,) '

If you want an empty string rather than NULL, write double quotes:

' ("",42,) '

Here the first field is a non-NULL empty string, the third is NULL.

(These constants are actually only a special case of the generic type constants discussed in Section 4.1.2.7. The constant is initially treated as a string and passed to the composite-type input conversion routine. An explicit type specification might be necessary to tell which type to convert the constant to.)

The ROW expression syntax can also be used to construct composite values. In most cases this is considerably simpler to use than the string-literal syntax since you don’t have to worry about multiple layers of quoting. We already used this method above:

ROW('fuzzy dice', 42, 1.99)
ROW('', 42, NULL)

The ROW keyword is actually optional as long as you have more than one field in the expression, so these can be simplified to:

('fuzzy dice', 42, 1.99)
('', 42, NULL)

The ROW expression syntax is discussed in more detail in Section 4.2.13.

8.16.3. Accessing Composite Types

To access a field of a composite column, one writes a dot and the field name, much like selecting a field from a table name. In fact, it’s so much like selecting from a table name that you often have to use parentheses to keep from confusing the parser. For example, you might try to select some subfields from our on_hand example table with something like:
SELECT item.name FROM on_hand WHERE item.price > 9.99;
This will not work since the name item is taken to be a table name, not a column name of on_hand, per SQL syntax rules. You must write it like this:

SELECT (item).name FROM on_hand WHERE (item).price > 9.99;
or if you need to use the table name as well (for instance in a multitable query), like this:

SELECT (on_hand.item).name FROM on_hand WHERE (on_hand.item).price > 9.99;
Now the parenthesized object is correctly interpreted as a reference to the item column, and then the subfield can be selected from it.

Similar syntactic issues apply whenever you select a field from a composite value. For instance, to select just one field from the result of a function that returns a composite value, you'd need to write something like:

SELECT (my_func(...)).field FROM ...
Without the extra parentheses, this will generate a syntax error.
The special field name * means “all fields”, as further explained in Section 8.16.5.

8.16.4. Modifying Composite Types
Here are some examples of the proper syntax for inserting and updating composite columns. First, inserting or updating a whole column:

INSERT INTO mytab (complex_col) VALUES((1.1,2.2));
UPDATE mytab SET complex_col = ROW(1.1,2.2) WHERE ...
The first example omits ROW, the second uses it; we could have done it either way.

We can update an individual subfield of a composite column:

UPDATE mytab SET complex_col.r = (complex_col).r + 1 WHERE ...
Notice here that we don’t need to (and indeed cannot) put parentheses around the column name appearing just after SET, but we do need parentheses when referencing the same column in the expression to the right of the equal sign.

And we can specify subfields as targets for INSERT, too:

INSERT INTO mytab (complex_col.r, complex_col.i) VALUES(1.1, 2.2);
Had we not supplied values for all the subfields of the column, the remaining subfields would have been filled with null values.

8.16.5. Using Composite Types in Queries
There are various special syntax rules and behaviors associated with composite types in queries. These rules provide useful shortcuts, but can be confusing if you don’t know the logic behind them.
In PostgreSQL, a reference to a table name (or alias) in a query is effectively a reference to the composite value of the table's current row. For example, if we had a table `inventory_item` as shown above, we could write:

```sql
SELECT c FROM inventory_item c;
```

This query produces a single composite-valued column, so we might get output like:

```
<table>
<thead>
<tr>
<th>c</th>
</tr>
</thead>
<tbody>
<tr>
<td>(&quot;fuzzy dice&quot;,42,1.99)</td>
</tr>
<tr>
<td>(1 row)</td>
</tr>
</tbody>
</table>
```

Note however that simple names are matched to column names before table names, so this example works only because there is no column named `c` in the query's tables.

The ordinary qualified-column-name syntax `table_name.column_name` can be understood as applying field selection to the composite value of the table's current row. (For efficiency reasons, it's not actually implemented that way.)

When we write

```sql
SELECT c.* FROM inventory_item c;
```

then, according to the SQL standard, we should get the contents of the table expanded into separate columns:

```
| name    | supplier_id | price |
|----------|-------------+-------|
| fuzzy dice | 42     | 1.99  |
| (1 row)  |
```

as if the query were

```sql
SELECT c.name, c.supplier_id, c.price FROM inventory_item c;
```

PostgreSQL will apply this expansion behavior to any composite-valued expression, although as shown above, you need to write parentheses around the value that `. *` is applied to whenever it's not a simple table name. For example, if `myfunc()` is a function returning a composite type with columns `a`, `b`, and `c`, then these two queries have the same result:

```sql
SELECT (myfunc(x)).* FROM some_table;
SELECT (myfunc(x)).a, (myfunc(x)).b, (myfunc(x)).c FROM some_table;
```

**Tip**

PostgreSQL handles column expansion by actually transforming the first form into the second. So, in this example, `myfunc()` would get invoked three times per row with either syntax. If it's an expensive function you may wish to avoid that, which you can do with a query like:

```sql
SELECT m.* FROM some_table, LATERAL myfunc(x) AS m;
```
Placing the function in a LATERAL FROM item keeps it from being invoked more than once per row. m.* is still expanded into m.a, m.b, m.c but now those variables are just references to the output of the FROM item. (The LATERAL keyword is optional here, but we show it to clarify that the function is getting x from some_table.)

The composite_value.* syntax results in column expansion of this kind when it appears at the top level of a SELECT output list, a RETURNING list in INSERT/UPDATE/DELETE, a VALUES clause, or a row constructor. In all other contexts (including when nested inside one of those constructs), attaching .* to a composite value does not change the value, since it means “all columns” and so the same composite value is produced again. For example, if somefunc() accepts a composite-valued argument, these queries are the same:

```
SELECT somefunc(c.*) FROM inventory_item c;
SELECT somefunc(c) FROM inventory_item c;
```

In both cases, the current row of inventory_item is passed to the function as a single composite-valued argument. Even though .* does nothing in such cases, using it is good style, since it makes clear that a composite value is intended. In particular, the parser will consider c in c.* to refer to a table name or alias, not to a column name, so that there is no ambiguity; whereas without .*, it is not clear whether c means a table name or a column name, and in fact the column-name interpretation will be preferred if there is a column named c.

Another example demonstrating these concepts is that all these queries mean the same thing:

```
SELECT * FROM inventory_item c ORDER BY c;
SELECT * FROM inventory_item c ORDER BY c.*;
SELECT * FROM inventory_item c ORDER BY ROW(c.*);
```

All of these ORDER BY clauses specify the row’s composite value, resulting in sorting the rows according to the rules described in Section 9.23.6. However, if inventory_item contained a column named c, the first case would be different from the others, as it would mean to sort by that column only. Given the column names previously shown, these queries are also equivalent to those above:

```
SELECT * FROM inventory_item c ORDER BY ROW(c.name, c.supplier_id, c.price);
SELECT * FROM inventory_item c ORDER BY (c.name, c.supplier_id, c.price);
```

(The last case uses a row constructor with the key word ROW omitted.)

Another special syntactical behavior associated with composite values is that we can use functional notation for extracting a field of a composite value. The simple way to explain this is that the notations field(table) and table.field are interchangeable. For example, these queries are equivalent:

```
SELECT c.name FROM inventory_item c WHERE c.price > 1000;
SELECT name(c) FROM inventory_item c WHERE price(c) > 1000;
```

Moreover, if we have a function that accepts a single argument of a composite type, we can call it with either notation. These queries are all equivalent:

```
SELECT somefunc(c) FROM inventory_item c;
SELECT somefunc(c.*) FROM inventory_item c;
```
SELECT c.somefunc FROM inventory_item c;

This equivalence between functional notation and field notation makes it possible to use functions on composite types to implement “computed fields”. An application using the last query above wouldn't need to be directly aware that somefunc isn't a real column of the table.

**Tip**

Because of this behavior, it's unwise to give a function that takes a single composite-type argument the same name as any of the fields of that composite type. If there is ambiguity, the field-name interpretation will be chosen if field-name syntax is used, while the function will be chosen if function-call syntax is used. However, PostgreSQL versions before 11 always chose the field-name interpretation, unless the syntax of the call required it to be a function call. One way to force the function interpretation in older versions is to schema-qualify the function name, that is, write `schema.func(compositevalue)`.

---

### 8.16.6. Composite Type Input and Output Syntax

The external text representation of a composite value consists of items that are interpreted according to the I/O conversion rules for the individual field types, plus decoration that indicates the composite structure. The decoration consists of parentheses ( ( and ) ) around the whole value, plus commas (,) between adjacent items. Whitespace outside the parentheses is ignored, but within the parentheses it is considered part of the field value, and might or might not be significant depending on the input conversion rules for the field data type. For example, in:

```
' (  42 )'
```

the whitespace will be ignored if the field type is integer, but not if it is text.

As shown previously, when writing a composite value you can write double quotes around any individual field value. You must do so if the field value would otherwise confuse the composite-value parser. In particular, fields containing parentheses, commas, double quotes, or backslashes must be double-quoted. To put a double quote or backslash in a quoted composite field value, precede it with a backslash. (Also, a pair of double quotes within a double-quoted field value is taken to represent a double quote character, analogously to the rules for single quotes in SQL literal strings.) Alternatively, you can avoid quoting and use backslash-escaping to protect all data characters that would otherwise be taken as composite syntax.

A completely empty field value (no characters at all between the commas or parentheses) represents a NULL. To write a value that is an empty string rather than NULL, write `""`.

The composite output routine will put double quotes around field values if they are empty strings or contain parentheses, commas, double quotes, backslashes, or white space. (Doing so for white space is not essential, but aids legibility.) Double quotes and backslashes embedded in field values will be doubled.

**Note**

Remember that what you write in an SQL command will first be interpreted as a string literal, and then as a composite. This doubles the number of backslashes you need (assuming escape string syntax is used). For example, to insert a text field containing a double quote and a backslash in a composite value, you'd need to write:

```
INSERT ... VALUES ('"\"\"');
```
Data Types

The string-literal processor removes one level of backslashes, so that what arrives at the composite-value parser looks like ("\""). In turn, the string fed to the text data type's input routine becomes \\.

(If we were working with a data type whose input routine also treated backslashes specially, bytea for example, we might need as many as eight backslashes in the command to get one backslash into the stored composite field.) Dollar quoting (see Section 4.1.2.4) can be used to avoid the need to double backslashes.

Tip

The ROW constructor syntax is usually easier to work with than the composite-literal syntax when writing composite values in SQL commands. In ROW, individual field values are written the same way they would be written when not members of a composite.

8.17. Range Types

Range types are data types representing a range of values of some element type (called the range's subtype). For instance, ranges of timestamp might be used to represent the ranges of time that a meeting room is reserved. In this case the data type is tsrange (short for “timestamp range”), and timestamp is the subtype. The subtype must have a total order so that it is well-defined whether element values are within, before, or after a range of values.

Range types are useful because they represent many element values in a single range value, and because concepts such as overlapping ranges can be expressed clearly. The use of time and date ranges for scheduling purposes is the clearest example; but price ranges, measurement ranges from an instrument, and so forth can also be useful.

8.17.1. Built-in Range Types

PostgreSQL comes with the following built-in range types:

- int4range — Range of integer
- int8range — Range of bigint
- numrange — Range of numeric
- tsrange — Range of timestamp without time zone
- tstzrange — Range of timestamp with time zone
- daterange — Range of date

In addition, you can define your own range types; see CREATE TYPE for more information.

8.17.2. Examples

```sql
CREATE TABLE reservation (room int, during tsrange);
INSERT INTO reservation VALUES
  (1108, '[2010-01-01 14:30, 2010-01-01 15:30]');

-- Containment
SELECT int4range(10, 20) @> 3;
```
-- Overlaps
SELECT numrange(11.1, 22.2) && numrange(20.0, 30.0);

-- Extract the upper bound
SELECT upper(int8range(15, 25));

-- Compute the intersection
SELECT int4range(10, 20) * int4range(15, 25);

-- Is the range empty?
SELECT isempty(numrange(1, 5));

See Table 9.50 and Table 9.51 for complete lists of operators and functions on range types.

8.17.3. Inclusive and Exclusive Bounds

Every non-empty range has two bounds, the lower bound and the upper bound. All points between these values are included in the range. An inclusive bound means that the boundary point itself is included in the range as well, while an exclusive bound means that the boundary point is not included in the range.

In the text form of a range, an inclusive lower bound is represented by "[" while an exclusive lower bound is represented by "]". Likewise, an inclusive upper bound is represented by "]", while an exclusive upper bound is represented by ")". (See Section 8.17.5 for more details.)

The functions lower_inc and upper_inc test the inclusivity of the lower and upper bounds of a range value, respectively.

8.17.4. Infinite (Unbounded) Ranges

The lower bound of a range can be omitted, meaning that all points less than the upper bound are included in the range. Likewise, if the upper bound of the range is omitted, then all points greater than the lower bound are included in the range. If both lower and upper bounds are omitted, all values of the element type are considered to be in the range.

This is equivalent to considering that the lower bound is “minus infinity”, or the upper bound is “plus infinity”, respectively. But note that these infinite values are never values of the range's element type, and can never be part of the range. (So there is no such thing as an inclusive infinite bound — if you try to write one, it will automatically be converted to an exclusive bound.)

Also, some element types have a notion of “infinity”, but that is just another value so far as the range type mechanisms are concerned. For example, in timestamp ranges, [today,] means the same thing as [today,). But [today, infinity] means something different from [today, infinity) — the latter excludes the special timestamp value infinity.

The functions lower_inf and upper_inf test for infinite lower and upper bounds of a range, respectively.

8.17.5. Range Input/Output

The input for a range value must follow one of the following patterns:

(lower-bound, upper-bound)
(lower-bound, upper-bound]
[lower-bound, upper-bound)
[lower-bound, upper-bound]
empty

The parentheses or brackets indicate whether the lower and upper bounds are exclusive or inclusive, as described previously. Notice that the final pattern is empty, which represents an empty range (a range that contains no points).

The lower-bound may be either a string that is valid input for the subtype, or empty to indicate no lower bound. Likewise, upper-bound may be either a string that is valid input for the subtype, or empty to indicate no upper bound.

Each bound value can be quoted using " (double quote) characters. This is necessary if the bound value contains parentheses, brackets, commas, double quotes, or backslashes, since these characters would otherwise be taken as part of the range syntax. To put a double quote or backslash in a quoted bound value, precede it with a backslash. (Also, a pair of double quotes within a double-quoted bound value is taken to represent a double quote character, analogously to the rules for single quotes in SQL literal strings.) Alternatively, you can avoid quoting and use backslash-escaping to protect all data characters that would otherwise be taken as range syntax. Also, to write a bound value that is an empty string, write ",", since writing nothing means an infinite bound.

Whitespace is allowed before and after the range value, but any whitespace between the parentheses or brackets is taken as part of the lower or upper bound value. (Depending on the element type, it might or might not be significant.)

Note

These rules are very similar to those for writing field values in composite-type literals. See Section 8.16.6 for additional commentary.

Examples:

-- includes 3, does not include 7, and does include all points in between
SELECT '[:3,7]'::int4range;

-- does not include either 3 or 7, but includes all points in between
SELECT '(3,7)'::int4range;

-- includes only the single point 4
SELECT '[:4,4]'::int4range;

-- includes no points (and will be normalized to 'empty')
SELECT '[:4,4]'::int4range;

8.17.6. Constructing Ranges

Each range type has a constructor function with the same name as the range type. Using the constructor function is frequently more convenient than writing a range literal constant, since it avoids the need for extra quoting of the bound values. The constructor function accepts two or three arguments. The two-argument form constructs a range in standard form (lower bound inclusive, upper bound exclusive), while the three-argument form constructs a range with bounds of the form specified by the third argument. The third argument must be one of the strings "()", "[]", "[)", or "\]". For example:

-- The full form is: lower bound, upper bound, and text argument indicating
-- inclusivity/exclusivity of bounds.
SELECT numrange(1.0, 14.0, '()');

-- If the third argument is omitted, '()' is assumed.
SELECT numrange(1.0, 14.0);

-- Although '()' is specified here, on display the value will be
-- canonical form, since int8range is a discrete range type (see
below).
SELECT int8range(1, 14, '()');

-- Using NULL for either bound causes the range to be unbounded on
that side.
SELECT numrange(NULL, 2.2);

8.17.7. Discrete Range Types

A discrete range is one whose element type has a well-defined “step”, such as integer or date.
In these types two elements can be said to be adjacent, when there are no valid values between them.
This contrasts with continuous ranges, where it’s always (or almost always) possible to identify other
element values between two given values. For example, a range over the numeric type is continuous,
as is a range over timestamp. (Even though timestamp has limited precision, and so could
theoretically be treated as discrete, it’s better to consider it continuous since the step size is normally
not of interest.)

Another way to think about a discrete range type is that there is a clear idea of a “next” or “previous”
value for each element value. Knowing that, it is possible to convert between inclusive and exclusive
representations of a range's bounds, by choosing the next or previous element value instead of the one
originally given. For example, in an integer range type [4,8) and (3,9) denote the same set of
values; but this would not be so for a range over numeric.

A discrete range type should have a canonicalization function that is aware of the desired step size for
the element type. The canonicalization function is charged with converting equivalent values of the
range type to have identical representations, in particular consistently inclusive or exclusive bounds.
If a canonicalization function is not specified, then ranges with different formatting will always be
treated as unequal, even though they might represent the same set of values in reality.

The built-in range types int4range, int8range, and daterange all use a canonical form that
includes the lower bound and excludes the upper bound; that is, {}. User-defined range types can use
other conventions, however.

8.17.8. Defining New Range Types

Users can define their own range types. The most common reason to do this is to use ranges over
subtypes not provided among the built-in range types. For example, to define a new range type of
subtype float8:

CREATE TYPE floatrange AS RANGE (
    subtype = float8,
    subtype_diff = float8mi
);

SELECT '[1.234, 5.678]'::floatrange;

Because float8 has no meaningful “step”, we do not define a canonicalization function in this ex-
ample.
Defining your own range type also allows you to specify a different subtype B-tree operator class or collation to use, so as to change the sort ordering that determines which values fall into a given range.

If the subtype is considered to have discrete rather than continuous values, the `CREATE TYPE` command should specify a canonical function. The canonicalization function takes an input range value, and must return an equivalent range value that may have different bounds and formatting. The canonical output for two ranges that represent the same set of values, for example the integer ranges \([1, 7]\) and \([1, 8)\), must be identical. It doesn't matter which representation you choose to be the canonical one, so long as two equivalent values with different formattings are always mapped to the same value with the same formatting. In addition to adjusting the inclusive/exclusive bounds format, a canonicalization function might round off boundary values, in case the desired step size is larger than what the subtype is capable of storing. For instance, a range type over `timestamp` could be defined to have a step size of an hour, in which case the canonicalization function would need to round off bounds that weren't a multiple of an hour, or perhaps throw an error instead.

In addition, any range type that is meant to be used with GiST or SP-GiST indexes should define a subtype difference, or `subtype_diff`, function. (The index will still work without `subtype_diff`, but it is likely to be considerably less efficient than if a difference function is provided.) The subtype difference function takes two input values of the subtype, and returns their difference (i.e., \(X\) minus \(Y\)) represented as a `float8` value. In our example above, the function `float8mi` that underlies the regular `float8` minus operator can be used; but for any other subtype, some type conversion would be necessary. Some creative thought about how to represent differences as numbers might be needed, too. To the greatest extent possible, the `subtype_diff` function should agree with the sort ordering implied by the selected operator class and collation; that is, its result should be positive whenever its first argument is greater than its second according to the sort ordering.

A less-oversimplified example of a `subtype_diff` function is:

```sql
CREATE FUNCTION time_subtype_diff(x time, y time) RETURNS float8 AS 'SELECT EXTRACT(EPOCH FROM (x - y))' LANGUAGE sql STRICT IMMUTABLE;
CREATE TYPE timerange AS RANGE (
    subtype = time,
    subtype_diff = time_subtype_diff
);
SELECT '([11:10, 23:00])'::timerange;
```

See `CREATE TYPE` for more information about creating range types.

### 8.17.9. Indexing

GiST and SP-GiST indexes can be created for table columns of range types. For instance, to create a GiST index:

```sql
CREATE INDEX reservation_idx ON reservation USING GIST (during);
```

A GiST or SP-GiST index can accelerate queries involving these range operators: `=`, `&`, `<@`, `@>`, `<<`, `>>`, `-|`, `<`, and `&>`. (see Table 9.50 for more information).

In addition, B-tree and hash indexes can be created for table columns of range types. For these index types, basically the only useful range operation is equality. There is a B-tree sort ordering defined for range values, with corresponding `<` and `>` operators, but the ordering is rather arbitrary and not usually useful in the real world. Range types' B-tree and hash support is primarily meant to allow sorting and hashing internally in queries, rather than creation of actual indexes.

### 8.17.10. Constraints on Ranges
While UNIQUE is a natural constraint for scalar values, it is usually unsuitable for range types. Instead, an exclusion constraint is often more appropriate (see CREATE TABLE ... CONSTRAINT ... EXCLUDE). Exclusion constraints allow the specification of constraints such as “non-overlapping” on a range type. For example:

```sql
CREATE TABLE reservation (  
during tsrange,  
EXCLUDE USING GIST (during WITH &&)  
);
```

That constraint will prevent any overlapping values from existing in the table at the same time:

```sql
INSERT INTO reservation VALUES  
('[2010-01-01 11:30, 2010-01-01 15:00]');
```

```sql
INSERT INTO reservation VALUES  
('[2010-01-01 14:45, 2010-01-01 15:45]');
ERROR: conflicting key value violates exclusion constraint  
"reservation_during_excl"  
DETAIL: Key (during)=("2010-01-01 14:45:00","2010-01-01 15:45:00") conflicts  
with existing key (during)=("2010-01-01 11:30:00","2010-01-01 15:00:00").
```

You can use the btree_gist extension to define exclusion constraints on plain scalar data types, which can then be combined with range exclusions for maximum flexibility. For example, after btree_gist is installed, the following constraint will reject overlapping ranges only if the meeting room numbers are equal:

```sql
CREATE EXTENSION btree_gist;  
CREATE TABLE room_reservation (  
room text,  
during tsrange,  
EXCLUDE USING GIST (room WITH =, during WITH &&)  
);
```

```sql
INSERT INTO room_reservation VALUES  
('123A', '[2010-01-01 14:00, 2010-01-01 15:00]');
```

```sql
INSERT INTO room_reservation VALUES  
('123A', '[2010-01-01 14:30, 2010-01-01 15:30]');
ERROR: conflicting key value violates exclusion constraint  
"room_reservation_room_during_excl"  
DETAIL: Key (room, during)=(123A, ["2010-01-01 14:30:00","2010-01-01 15:30:00"]) conflicts  
with existing key (room, during)=(123A, ["2010-01-01 14:00:00","2010-01-01 15:00:00"]).
```

```sql
INSERT INTO room_reservation VALUES  
('123B', '[2010-01-01 14:30, 2010-01-01 15:30]');
```

## 8.18. Domain Types
A *domain* is a user-defined data type that is based on another *underlying type*. Optionally, it can have constraints that restrict its valid values to a subset of what the underlying type would allow. Otherwise it behaves like the underlying type — for example, any operator or function that can be applied to the underlying type will work on the domain type. The underlying type can be any built-in or user-defined base type, enum type, array type, composite type, range type, or another domain.

For example, we could create a domain over integers that accepts only positive integers:

```
CREATE DOMAIN posint AS integer CHECK (VALUE > 0);
CREATE TABLE mytable (id posint);
INSERT INTO mytable VALUES(1);   -- works
INSERT INTO mytable VALUES(-1);  -- fails
```

When an operator or function of the underlying type is applied to a domain value, the domain is automatically down-cast to the underlying type. Thus, for example, the result of `mytable.id - 1` is considered to be of type `integer` not `posint`. We could write `(mytable.id - 1)::posint` to cast the result back to `posint`, causing the domain's constraints to be rechecked. In this case, that would result in an error if the expression had been applied to an `id` value of 1. Assigning a value of the underlying type to a field or variable of the domain type is allowed without writing an explicit cast, but the domain's constraints will be checked.

For additional information see CREATE DOMAIN.

### 8.19. Object Identifier Types

Object identifiers (OIDs) are used internally by PostgreSQL as primary keys for various system tables. OIDs are not added to user-created tables, unless `WITH OIDS` is specified when the table is created, or the `default_with_oids` configuration variable is enabled. Type `oid` represents an object identifier. There are also several alias types for `oid`: `regproc`, `regprocedure`, `regoper`, `regoperator`, `regclass`, `regtype`, `regrole`, `regnamespace`, `regconfig`, and `regdictionary`. Table 8.24 shows an overview.

The `oid` type is currently implemented as an unsigned four-byte integer. Therefore, it is not large enough to provide database-wide uniqueness in large databases, or even in large individual tables. So, using a user-created table's OID column as a primary key is discouraged. OIDs are best used only for references to system tables.

The `oid` type itself has few operations beyond comparison. It can be cast to integer, however, and then manipulated using the standard integer operators. (Beware of possible signed-versus-unsigned confusion if you do this.)

The OID alias types have no operations of their own except for specialized input and output routines. These routines are able to accept and display symbolic names for system objects, rather than the raw numeric value that type `oid` would use. The alias types allow simplified lookup of OID values for objects. For example, to examine the `pg_attribute` rows related to a table `mytable`, one could write:

```
SELECT * FROM pg_attribute WHERE attrelid = 'mytable'::regclass;
```

rather than:

```
SELECT * FROM pg_attribute
  WHERE attrelid = (SELECT oid FROM pg_class WHERE relname = 'mytable');
```

While that doesn't look all that bad by itself, it's still oversimplified. A far more complicated sub-select would be needed to select the right OID if there are multiple tables named `mytable` in different
Data Types

schemas. The regclass input converter handles the table lookup according to the schema path setting, and so it does the "right thing" automatically. Similarly, casting a table's OID to regclass is handy for symbolic display of a numeric OID.

Table 8.24. Object Identifier Types

<table>
<thead>
<tr>
<th>Name</th>
<th>References</th>
<th>Description</th>
<th>Value Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>any</td>
<td>numeric object identi-</td>
<td>564182</td>
</tr>
<tr>
<td>regproc</td>
<td>pg_proc</td>
<td>function name</td>
<td>sum</td>
</tr>
<tr>
<td>regprocedure</td>
<td>pg_proc</td>
<td>function with argument</td>
<td>sum(int4)</td>
</tr>
<tr>
<td>regoper</td>
<td>pg_operator</td>
<td>operator name</td>
<td>+</td>
</tr>
<tr>
<td>regoperator</td>
<td>pg_operator</td>
<td>operator with argument</td>
<td>*(integer,integer)</td>
</tr>
<tr>
<td>regclass</td>
<td>pg_class</td>
<td>relation name</td>
<td>pg_type</td>
</tr>
<tr>
<td>regtype</td>
<td>pg_type</td>
<td>data type name</td>
<td>integer</td>
</tr>
<tr>
<td>regrole</td>
<td>pg_authid</td>
<td>role name</td>
<td>smithee</td>
</tr>
<tr>
<td>regnamespace</td>
<td>pg_namespace</td>
<td>namespace name</td>
<td>pg_catalog</td>
</tr>
<tr>
<td>regconfig</td>
<td>pg_ts_config</td>
<td>text search configura-</td>
<td>english</td>
</tr>
<tr>
<td>regdictionary</td>
<td>pg_ts_dict</td>
<td>text search dictionary</td>
<td>simple</td>
</tr>
</tbody>
</table>

All of the OID alias types for objects grouped by namespace accept schema-qualified names, and will display schema-qualified names on output if the object would not be found in the current search path without being qualified. The regproc and regoperator alias types will only accept input names that are unique (not overloaded), so they are of limited use; for most uses regprocedure or regoperator are more appropriate. For regoperator, unary operators are identified by writing NONE for the unused operand.

An additional property of most of the OID alias types is the creation of dependencies. If a constant of one of these types appears in a stored expression (such as a column default expression or view), it creates a dependency on the referenced object. For example, if a column has a default expression

\[ \text{nextval('my_seq'::regclass)} \]

PostgreSQL understands that the default expression depends on the sequence my_seq; the system will not let the sequence be dropped without first removing the default expression. regrole is the only exception for the property. Constants of this type are not allowed in such expressions.

Note

The OID alias types do not completely follow transaction isolation rules. The planner also treats them as simple constants, which may result in sub-optimal planning.

Another identifier type used by the system is xid, or transaction (abbreviated xact) identifier. This is the data type of the system columns xmin and xmax. Transaction identifiers are 32-bit quantities.

A third identifier type used by the system is cid, or command identifier. This is the data type of the system columns cmin and cmax. Command identifiers are also 32-bit quantities.

A final identifier type used by the system is tid, or tuple identifier (row identifier). This is the data type of the system column ctid. A tuple ID is a pair (block number, tuple index within block) that identifies the physical location of the row within its table.
8.20. pg_lsn Type

The `pg_lsn` data type can be used to store LSN (Log Sequence Number) data which is a pointer to a location in the WAL. This type is a representation of `XLogRecPtr` and an internal system type of PostgreSQL.

Internally, an LSN is a 64-bit integer, representing a byte position in the write-ahead log stream. It is printed as two hexadecimal numbers of up to 8 digits each, separated by a slash; for example, `16/B374D848`. The `pg_lsn` type supports the standard comparison operators, like `=` and `>`. Two LSNs can be subtracted using the `-` operator; the result is the number of bytes separating those write-ahead log locations.

8.21. Pseudo-Types

The PostgreSQL type system contains a number of special-purpose entries that are collectively called pseudo-types. A pseudo-type cannot be used as a column data type, but it can be used to declare a function's argument or result type. Each of the available pseudo-types is useful in situations where a function's behavior does not correspond to simply taking or returning a value of a specific SQL data type. Table 8.25 lists the existing pseudo-types.

### Table 8.25. Pseudo-Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>any</td>
<td>Indicates that a function accepts any input data type.</td>
</tr>
<tr>
<td>anyelement</td>
<td>Indicates that a function accepts any data type (see Section 38.2.5).</td>
</tr>
<tr>
<td>anyarray</td>
<td>Indicates that a function accepts any array data type (see Section 38.2.5).</td>
</tr>
<tr>
<td>anynonarray</td>
<td>Indicates that a function accepts any non-array data type (see Section 38.2.5).</td>
</tr>
<tr>
<td>anyenum</td>
<td>Indicates that a function accepts any enum data type (see Section 38.2.5 and Section 8.7).</td>
</tr>
<tr>
<td>anyrange</td>
<td>Indicates that a function accepts any range data type (see Section 38.2.5 and Section 8.17).</td>
</tr>
<tr>
<td>cstring</td>
<td>Indicates that a function accepts or returns a null-terminated C string.</td>
</tr>
<tr>
<td>internal</td>
<td>Indicates that a function accepts or returns a server-internal data type.</td>
</tr>
<tr>
<td>language_handler</td>
<td>A procedural language call handler is declared to return <code>language_handler</code>.</td>
</tr>
<tr>
<td>fdw_handler</td>
<td>A foreign-data wrapper handler is declared to return <code>fdw_handler</code>.</td>
</tr>
<tr>
<td>index_am_handler</td>
<td>An index access method handler is declared to return <code>index_am_handler</code>.</td>
</tr>
<tr>
<td>tsm_handler</td>
<td>A tablesample method handler is declared to return <code>tsm_handler</code>.</td>
</tr>
<tr>
<td>record</td>
<td>Identifies a function taking or returning an unspecified row type.</td>
</tr>
</tbody>
</table>
Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger</td>
<td>A trigger function is declared to return trigger.</td>
</tr>
<tr>
<td>event_trigger</td>
<td>An event trigger function is declared to return event_trigger.</td>
</tr>
<tr>
<td>pg ddl_command</td>
<td>Identifies a representation of DDL commands that is available to event triggers.</td>
</tr>
<tr>
<td>void</td>
<td>Indicates that a function returns no value.</td>
</tr>
<tr>
<td>unknown</td>
<td>Identifies a not-yet-resolved type, e.g. of an undecorated string literal.</td>
</tr>
<tr>
<td>opaque</td>
<td>An obsolete type name that formerly served many of the above purposes.</td>
</tr>
</tbody>
</table>

Functions coded in C (whether built-in or dynamically loaded) can be declared to accept or return any of these pseudo data types. It is up to the function author to ensure that the function will behave safely when a pseudo-type is used as an argument type.

Functions coded in procedural languages can use pseudo-types only as allowed by their implementation languages. At present most procedural languages forbid use of a pseudo-type as an argument type, and allow only void and record as a result type (plus trigger or event_trigger when the function is used as a trigger or event trigger). Some also support polymorphic functions using the types anyelement, anyarray, anynonarray, anyenum, and anyrange.

The internal pseudo-type is used to declare functions that are meant only to be called internally by the database system, and not by direct invocation in an SQL query. If a function has at least one internal-type argument then it cannot be called from SQL. To preserve the type safety of this restriction it is important to follow this coding rule: do not create any function that is declared to return internal unless it has at least one internal argument.
Chapter 9. Functions and Operators

PostgreSQL provides a large number of functions and operators for the built-in data types. Users can also define their own functions and operators, as described in Part V. The psql commands \df and \do can be used to list all available functions and operators, respectively.

If you are concerned about portability then note that most of the functions and operators described in this chapter, with the exception of the most trivial arithmetic and comparison operators and some explicitly marked functions, are not specified by the SQL standard. Some of this extended functionality is present in other SQL database management systems, and in many cases this functionality is compatible and consistent between the various implementations. This chapter is also not exhaustive; additional functions appear in relevant sections of the manual.

9.1. Logical Operators

The usual logical operators are available:

AND
OR
NOT

SQL uses a three-valued logic system with true, false, and null, which represents “unknown”. Observe the following truth tables:

<table>
<thead>
<tr>
<th>a</th>
<th>b</th>
<th>a AND b</th>
<th>a OR b</th>
</tr>
</thead>
<tbody>
<tr>
<td>TRUE</td>
<td>TRUE</td>
<td>TRUE</td>
<td>TRUE</td>
</tr>
<tr>
<td>TRUE</td>
<td>FALSE</td>
<td>FALSE</td>
<td>TRUE</td>
</tr>
<tr>
<td>TRUE</td>
<td>NULL</td>
<td>NULL</td>
<td>TRUE</td>
</tr>
<tr>
<td>FALSE</td>
<td>FALSE</td>
<td>FALSE</td>
<td>FALSE</td>
</tr>
<tr>
<td>FALSE</td>
<td>NULL</td>
<td>FALSE</td>
<td>NULL</td>
</tr>
<tr>
<td>NULL</td>
<td>NULL</td>
<td>NULL</td>
<td>NULL</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>a</th>
<th>NOT a</th>
</tr>
</thead>
<tbody>
<tr>
<td>TRUE</td>
<td>FALSE</td>
</tr>
<tr>
<td>FALSE</td>
<td>TRUE</td>
</tr>
<tr>
<td>NULL</td>
<td>NULL</td>
</tr>
</tbody>
</table>

The operators AND and OR are commutative, that is, you can switch the left and right operand without affecting the result. But see Section 4.2.14 for more information about the order of evaluation of subexpressions.

9.2. Comparison Functions and Operators

The usual comparison operators are available, as shown in Table 9.1.

**Table 9.1. Comparison Operators**

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;</td>
<td>less than</td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal to</td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater than or equal to</td>
</tr>
</tbody>
</table>
Comparison operators are available for all relevant data types. All comparison operators are binary operators that return values of type boolean; expressions like 1 < 2 < 3 are not valid (because there is no < operator to compare a Boolean value with 3).

There are also some comparison predicates, as shown in Table 9.2. These behave much like operators, but have special syntax mandated by the SQL standard.

### Table 9.2. Comparison Predicates

<table>
<thead>
<tr>
<th>Predicate</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>a BETWEEN x AND y</td>
<td>between</td>
</tr>
<tr>
<td>a NOT BETWEEN x AND y</td>
<td>not between</td>
</tr>
<tr>
<td>a BETWEEN SYMMETRIC x AND y</td>
<td>between, after sorting the comparison values</td>
</tr>
<tr>
<td>a NOT BETWEEN SYMMETRIC x AND y</td>
<td>not between, after sorting the comparison values</td>
</tr>
<tr>
<td>a IS DISTINCT FROM b</td>
<td>not equal, treating null like an ordinary value</td>
</tr>
<tr>
<td>a IS NOT DISTINCT FROM b</td>
<td>equal, treating null like an ordinary value</td>
</tr>
<tr>
<td>expression IS NULL</td>
<td>is null</td>
</tr>
<tr>
<td>expression IS NOT NULL</td>
<td>is not null</td>
</tr>
<tr>
<td>expression ISNULL</td>
<td>is null (nonstandard syntax)</td>
</tr>
<tr>
<td>expression NOTNULL</td>
<td>is not null (nonstandard syntax)</td>
</tr>
<tr>
<td>boolean_expression IS TRUE</td>
<td>is true</td>
</tr>
<tr>
<td>boolean_expression IS NOT TRUE</td>
<td>is false or unknown</td>
</tr>
<tr>
<td>boolean_expression IS FALSE</td>
<td>is false</td>
</tr>
<tr>
<td>boolean_expression IS NOT FALSE</td>
<td>is true or unknown</td>
</tr>
<tr>
<td>boolean_expression IS UNKNOWN</td>
<td>is unknown</td>
</tr>
<tr>
<td>boolean_expression IS NOT UNKNOWN</td>
<td>is true or false</td>
</tr>
</tbody>
</table>

The BETWEEN predicate simplifies range tests:

```
a BETWEEN x AND y
```

is equivalent to

```
a >= x AND a <= y
```

Notice that BETWEEN treats the endpoint values as included in the range. NOT BETWEEN does the opposite comparison:

```
a NOT BETWEEN x AND y
```
is equivalent to

\[ a < x \text{ OR } a > y \]

\( \text{BETWEEN SYMMETRIC} \) is like \( \text{BETWEEN} \) except there is no requirement that the argument to the left of \( \text{AND} \) be less than or equal to the argument on the right. If it is not, those two arguments are automatically swapped, so that a nonempty range is always implied.

Ordinary comparison operators yield null (signifying “unknown”), not true or false, when either input is null. For example, \( ? = \text{NULL} \) yields null, as does \( ? <> \text{NULL} \). When this behavior is not suitable, use the \( \text{IS [ NOT ] DISTINCT FROM} \) predicates:

\[ a \text{ IS DISTINCT FROM } b \]
\[ a \text{ IS NOT DISTINCT FROM } b \]

For non-null inputs, \( \text{IS DISTINCT FROM} \) is the same as the \(<>\) operator. However, if both inputs are null it returns false, and if only one input is null it returns true. Similarly, \( \text{IS NOT DISTINCT FROM} \) is identical to \( = \) for non-null inputs, but it returns true when both inputs are null, and false when only one input is null. Thus, these predicates effectively act as though null were a normal data value, rather than “unknown”.

To check whether a value is or is not null, use the predicates:

\[ \text{expression IS NULL} \]
\[ \text{expression IS NOT NULL} \]

or the equivalent, but nonstandard, predicates:

\[ \text{expression ISNULL} \]
\[ \text{expression NOTNULL} \]

Do not write \( \text{expression} = \text{NULL} \) because \( \text{NULL} \) is not “equal to” \( \text{NULL} \). (The null value represents an unknown value, and it is not known whether two unknown values are equal.)

### Tip

Some applications might expect that \( \text{expression} = \text{NULL} \) returns true if \( \text{expression} \) evaluates to the null value. It is highly recommended that these applications be modified to comply with the SQL standard. However, if that cannot be done the \text{transform_null_equals} configuration variable is available. If it is enabled, PostgreSQL will convert \( x = \text{NULL} \) clauses to \( x \text{ IS NULL} \).

If the \( \text{expression} \) is row-valued, then \( \text{IS NULL} \) is true when the row expression itself is null or when all the row’s fields are null, while \( \text{IS NOT NULL} \) is true when the row expression itself is non-null and all the row’s fields are non-null. Because of this behavior, \( \text{IS NULL} \) and \( \text{IS NOT NULL} \) do not always return inverse results for row-valued expressions; in particular, a row-valued expression that contains both null and non-null fields will return false for both tests. In some cases, it may be preferable to write \( \text{row IS DISTINCT FROM NULL} \) or \( \text{row IS NOT DISTINCT FROM NULL} \), which will simply check whether the overall row value is null without any additional tests on the row fields.

Boolean values can also be tested using the predicates

\[ \text{boolean_expression IS TRUE} \]
Functions and Operators

A null input is treated as the logical value "unknown". Notice that IS UNKNOWN and IS NOT UNKNOWN are effectively the same as IS NULL and IS NOT NULL, respectively, except that the input expression must be of Boolean type.

Some comparison-related functions are also available, as shown in Table 9.3.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>num_nonnulls(ANY)</td>
<td>returns the number of non-null arguments</td>
<td>num_nonnulls(1, NULL, 2)</td>
<td>2</td>
</tr>
<tr>
<td>num_nulls(VARIADIC &quot;any&quot;)</td>
<td>returns the number of null arguments</td>
<td>num_nulls(1, NULL, 2)</td>
<td>1</td>
</tr>
</tbody>
</table>

9.3. Mathematical Functions and Operators

Mathematical operators are provided for many PostgreSQL types. For types without standard mathematical conventions (e.g., date/time types) we describe the actual behavior in subsequent sections.

Table 9.4 shows the available mathematical operators.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>addition</td>
<td>2 + 3</td>
<td>5</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>2 - 3</td>
<td>-1</td>
</tr>
<tr>
<td>*</td>
<td>multiplication</td>
<td>2 * 3</td>
<td>6</td>
</tr>
<tr>
<td>/</td>
<td>division</td>
<td>4 / 2</td>
<td>2</td>
</tr>
<tr>
<td>%</td>
<td>modulo</td>
<td>5 % 4</td>
<td>1</td>
</tr>
<tr>
<td>^</td>
<td>exponentiation</td>
<td>2.0 ^ 3.0</td>
<td>8</td>
</tr>
<tr>
<td>/\</td>
<td>square root</td>
<td>/\ 25.0</td>
<td>5</td>
</tr>
<tr>
<td>//</td>
<td>cube root</td>
<td>// 27.0</td>
<td>3</td>
</tr>
<tr>
<td>!</td>
<td>factorial</td>
<td>5 !</td>
<td>120</td>
</tr>
<tr>
<td>!!</td>
<td>factorial (prefix operator)</td>
<td>!! 5</td>
<td>120</td>
</tr>
<tr>
<td>@</td>
<td>absolute value</td>
<td>@ -5.0</td>
<td>5</td>
</tr>
<tr>
<td>&amp;</td>
<td>bitwise AND</td>
<td>91 &amp; 15</td>
<td>11</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bitwise OR</td>
<td>32</td>
</tr>
<tr>
<td>#</td>
<td>bitwise XOR</td>
<td>17 # 5</td>
<td>20</td>
</tr>
<tr>
<td>~</td>
<td>bitwise NOT</td>
<td>~1</td>
<td>-2</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>bitwise shift left</td>
<td>1 &lt;&lt; 4</td>
<td>16</td>
</tr>
</tbody>
</table>
The bitwise operators work only on integral data types, whereas the others are available for all numeric data types. The bitwise operators are also available for the bit string types `bit` and `bit varying`, as shown in Table 9.13.

Table 9.5 shows the available mathematical functions. In the table, `dp` indicates double precision. Many of these functions are provided in multiple forms with different argument types. Except where noted, any given form of a function returns the same data type as its argument. The functions working with double precision data are mostly implemented on top of the host system’s C library; accuracy and behavior in boundary cases can therefore vary depending on the host system.

### Table 9.5. Mathematical Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>abs(x)</code></td>
<td>(same as input)</td>
<td>absolute value</td>
<td><code>abs(-17.4)</code></td>
<td>17.4</td>
</tr>
<tr>
<td><code>cbrt(dp)</code></td>
<td><code>dp</code></td>
<td>cube root</td>
<td><code>cbrt(27.0)</code></td>
<td>3</td>
</tr>
<tr>
<td><code>ceil(dp)</code></td>
<td><code>dp</code></td>
<td>nearest integer greater than or equal to argument</td>
<td><code>ceil(-42.8)</code></td>
<td>-42</td>
</tr>
<tr>
<td><code>ceiling(dp)</code></td>
<td><code>dp</code></td>
<td>nearest integer greater than or equal to argument (same as <code>ceil</code>)</td>
<td><code>ceiling(-95.3)</code></td>
<td>-95</td>
</tr>
<tr>
<td><code>degrees(dp)</code></td>
<td><code>dp</code></td>
<td>radians to degrees</td>
<td><code>degrees(0.5)</code></td>
<td>28.6478897565412</td>
</tr>
<tr>
<td><code>div(y numeric, x numeric)</code></td>
<td><code>numeric</code></td>
<td>integer quotient of <code>y/x</code></td>
<td><code>div(9,4)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>exp(dp or numeric)</code></td>
<td><code>dp</code></td>
<td>exponential</td>
<td><code>exp(1.0)</code></td>
<td>2.71828182845905</td>
</tr>
<tr>
<td><code>floor(dp or numeric)</code></td>
<td><code>dp</code></td>
<td>nearest integer less than or equal to argument</td>
<td><code>floor(-42.8)</code></td>
<td>-43</td>
</tr>
<tr>
<td><code>ln(dp or numeric)</code></td>
<td><code>dp</code></td>
<td>natural logarithm</td>
<td><code>ln(2.0)</code></td>
<td>0.693147180559945</td>
</tr>
<tr>
<td><code>log(dp or numeric)</code></td>
<td><code>dp</code></td>
<td>base 10 logarithm</td>
<td><code>log(100.0)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>log(b numeric, x numeric)</code></td>
<td><code>numeric</code></td>
<td>logarithm to base <code>b</code></td>
<td><code>log(2.0, 64.0)</code></td>
<td>6.00000000000</td>
</tr>
<tr>
<td><code>mod(y, x)</code></td>
<td>(same as argument types)</td>
<td>remainder of <code>y/x</code></td>
<td><code>mod(9,4)</code></td>
<td>1</td>
</tr>
<tr>
<td><code>pi()</code></td>
<td><code>dp</code></td>
<td>“#” constant</td>
<td><code>pi()</code></td>
<td>3.14159265358979</td>
</tr>
<tr>
<td><code>power(a dp, b dp)</code></td>
<td><code>dp</code></td>
<td><code>a</code> raised to the power of <code>b</code></td>
<td><code>power(9.0, 3.0)</code></td>
<td>729</td>
</tr>
<tr>
<td><code>power(a numeric, b numeric)</code></td>
<td><code>numeric</code></td>
<td><code>a</code> raised to the power of <code>b</code></td>
<td><code>power(9.0, 3.0)</code></td>
<td>729</td>
</tr>
<tr>
<td><code>radians(dp)</code></td>
<td><code>dp</code></td>
<td>degrees to radians</td>
<td><code>radians(45.0)</code></td>
<td>0.785398163397448</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------------------------</td>
<td>----------------------------------</td>
<td>-------------------------------------------------------</td>
<td>-----------</td>
<td>------------</td>
</tr>
<tr>
<td><code>round(dp or numeric)</code></td>
<td>(same as input)</td>
<td>round to nearest integer</td>
<td><code>round(42.4)</code></td>
<td>42</td>
</tr>
<tr>
<td><code>round(v numeric, s int)</code></td>
<td>numeric</td>
<td>round to <code>s</code> decimal places</td>
<td><code>round(42.4382,42.44)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>scale(numeric)</code></td>
<td></td>
<td>scale of the argument (the number of decimal digits in the fractional part)</td>
<td><code>scale(8.41)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>sign(dp or numeric)</code></td>
<td>(same as input)</td>
<td>sign of the argument (-1, 0, +1)</td>
<td><code>sign(-8.4)</code></td>
<td>-1</td>
</tr>
<tr>
<td><code>sqrt(dp or numeric)</code></td>
<td>(same as input)</td>
<td>square root</td>
<td><code>sqrt(2.0)</code></td>
<td>1.4142135623731</td>
</tr>
<tr>
<td><code>trunc(dp or numeric)</code></td>
<td>(same as input)</td>
<td>truncate toward zero</td>
<td><code>trunc(42.8)</code></td>
<td>42</td>
</tr>
<tr>
<td><code>trunc(v numeric, s int)</code></td>
<td>numeric</td>
<td>truncate to <code>s</code> decimal places</td>
<td><code>trunc(42.4382,42.43)</code></td>
<td>2</td>
</tr>
</tbody>
</table>

```
width_bucket(operand dp, b1 dp, b2 dp, count int)  return the bucket number to which operand would be assigned in a histogram having count equal-width buckets spanning the range b1 to b2; returns 0 or count+1 for an input outside the range

width_bucket(operand numeric, b1 numeric, b2 numeric, count int)  return the bucket number to which operand would be assigned in a histogram having count equal-width buckets spanning the range b1 to b2; returns 0 or count+1 for an input outside the range

width_bucket(operand anyelement, thresholds anyarray)  return the bucket number to which operand would be assigned in an array listing the lower bounds of the buckets; returns 0 for an input less than the first lower bound; the thresholds
```

203
Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>array</td>
<td>must be sort-ed, smallest first, or unexpected results will be obtained</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 9.6 shows functions for generating random numbers.

**Table 9.6. Random Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>random()</td>
<td>dp</td>
<td>random value in the range 0.0 &lt;= x &lt; 1.0</td>
</tr>
<tr>
<td>setseed(dp)</td>
<td>void</td>
<td>set seed for subsequent random() calls (value between -1.0 and 1.0, inclusive)</td>
</tr>
</tbody>
</table>

The characteristics of the values returned by random() depend on the system implementation. It is not suitable for cryptographic applications; see pgcrypto module for an alternative.

Finally, Table 9.7 shows the available trigonometric functions. All trigonometric functions take arguments and return values of type double precision. Each of the trigonometric functions comes in two variants, one that measures angles in radians and one that measures angles in degrees.

**Table 9.7. Trigonometric Functions**

<table>
<thead>
<tr>
<th>Function (radians)</th>
<th>Function (degrees)</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>acos(x)</td>
<td>acosd(x)</td>
<td>inverse cosine</td>
</tr>
<tr>
<td>asin(x)</td>
<td>asind(x)</td>
<td>inverse sine</td>
</tr>
<tr>
<td>atan(x)</td>
<td>atand(x)</td>
<td>inverse tangent</td>
</tr>
<tr>
<td>atan2(y, x)</td>
<td>atan2d(y, x)</td>
<td>inverse tangent of y/x</td>
</tr>
<tr>
<td>cos(x)</td>
<td>cosd(x)</td>
<td>cosine</td>
</tr>
<tr>
<td>cot(x)</td>
<td>cotd(x)</td>
<td>cotangent</td>
</tr>
<tr>
<td>sin(x)</td>
<td>sind(x)</td>
<td>sine</td>
</tr>
<tr>
<td>tan(x)</td>
<td>tand(x)</td>
<td>tangent</td>
</tr>
</tbody>
</table>

**Note**

Another way to work with angles measured in degrees is to use the unit transformation functions radians() and degrees() shown earlier. However, using the degree-based trigonometric functions is preferred, as that way avoids round-off error for special cases such as sind(30).

9.4. String Functions and Operators

This section describes functions and operators for examining and manipulating string values. Strings in this context include values of the types character, character varying, and text. Unless otherwise noted, all of the functions listed below work on all of these types, but be wary of potential effects of automatic space-padding when using the character type. Some functions also exist natively for the bit-string types.
SQL defines some string functions that use key words, rather than commas, to separate arguments. Details are in Table 9.8. PostgreSQL also provides versions of these functions that use the regular function invocation syntax (see Table 9.9).

**Note**

Before PostgreSQL 8.3, these functions would silently accept values of several non-string data types as well, due to the presence of implicit coercions from those data types to `text`. Those coercions have been removed because they frequently caused surprising behaviors. However, the string concatenation operator (`||`) still accepts non-string input, so long as at least one input is of a string type, as shown in Table 9.8. For other cases, insert an explicit coercion to `text` if you need to duplicate the previous behavior.

### Table 9.8. SQL String Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>string</code></td>
<td><code>text</code></td>
<td>String concatenation</td>
<td>'Post' `</td>
<td>greSQL'</td>
</tr>
<tr>
<td><code>string</code></td>
<td><code>text</code></td>
<td>String concatenation with one non-string input</td>
<td>'Value: ' 42</td>
<td>Value: 42</td>
</tr>
<tr>
<td><code>bit_length(string)</code></td>
<td><code>int</code></td>
<td>Number of bits in string</td>
<td><code>bit_length('jose')</code></td>
<td>32</td>
</tr>
<tr>
<td><code>char_length(string)</code></td>
<td><code>int</code></td>
<td>Number of characters in string</td>
<td><code>char_length('jose')</code></td>
<td>4</td>
</tr>
<tr>
<td><code>lower(string)</code></td>
<td><code>text</code></td>
<td>Convert string to lower case</td>
<td><code>lower('TOM')</code></td>
<td>tom</td>
</tr>
<tr>
<td><code>octet_length(string)</code></td>
<td><code>int</code></td>
<td>Number of bytes in string</td>
<td><code>octet_length('jose')</code></td>
<td></td>
</tr>
<tr>
<td><code>overlay(string placing string from int [for int])</code></td>
<td><code>text</code></td>
<td>Replace substring</td>
<td><code>overlay('Txxxxas' placing 'hom' from 2 for 4)</code></td>
<td>Thomas</td>
</tr>
<tr>
<td><code>position(sub-string in string)</code></td>
<td><code>int</code></td>
<td>Location of specified substring</td>
<td><code>position('om' in 'Thomas')</code></td>
<td>3</td>
</tr>
<tr>
<td><code>substring(string [from int] [for int])</code></td>
<td><code>text</code></td>
<td>Extract substring matching POSIX regular expression</td>
<td><code>substring('Thomas' from 2 for 3)</code></td>
<td>hom</td>
</tr>
<tr>
<td><code>substring(string from pattern)</code></td>
<td><code>text</code></td>
<td>Extract substring</td>
<td><code>substring('Thomas' from '....$')</code></td>
<td>mas</td>
</tr>
</tbody>
</table>

See Section 9.7 for more information.
### Table 9.9. Other String Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>sub-string</td>
<td>text</td>
<td>Extract substring matching SQL regular expression from 'Thomas' for escape</td>
<td>sub-string('Thomas' from pattern for escape)</td>
<td></td>
</tr>
<tr>
<td>trim([lead-ing</td>
<td>trail-ing</td>
<td>both] [characters] from string)</td>
<td>Remove longest string containing only characters from characters (a space by default) from the start, end, or both ends (both is the default) of string</td>
<td>trim(both 'xyz' from 'yxTomxx')</td>
</tr>
<tr>
<td>btrim(string text [text [characters text]})</td>
<td>text</td>
<td>Remove the longest string consisting only of characters in characters (a space by default) from the start and end of string</td>
<td>btrim('xyxtrimyyx', 'xyz')</td>
<td></td>
</tr>
</tbody>
</table>

Additional string manipulation functions are available and are listed in Table 9.9. Some of them are used internally to implement the SQL-standard string functions listed in Table 9.8.
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>chr(int)</code></td>
<td>text</td>
<td>Character with the given code. For UTF8 the argument is treated as a Unicode code point. For other multibyte encodings the argument must designate an ASCII character. The NULL (0) character is not allowed because text data types cannot store such bytes.</td>
<td><code>chr(65)</code></td>
<td>A</td>
</tr>
<tr>
<td><code>concat(str text</code></td>
<td>text</td>
<td>Concatenate the text representations of all the arguments. NULL arguments are ignored.</td>
<td><code>concat('abcde', 2, NULL, 22)</code></td>
<td><code>abcde222</code></td>
</tr>
<tr>
<td><code>concat_ws(sep, str text, str, ...)</code></td>
<td>text</td>
<td>Concatenate all but the first argument with separators. The first argument is used as the separator string. NULL arguments are ignored.</td>
<td><code>concat_ws(',', 'abcde', 2, NULL, 22)</code></td>
<td><code>abcde,2,22</code></td>
</tr>
<tr>
<td><code>convert(string bytea, src_encoding name, dest_encoding name)</code></td>
<td>bytea</td>
<td>Convert string to the data-base encoding. The original encoding is specified by <code>src_encoding</code>. The string must be valid in this encoding. Conversions can be defined by <code>CREATE CONVERSION</code>. Also there are some predefined conversions. See Table 9.10 for available conversions.</td>
<td><code>convert('text_in_utf8', 'UTF8', 'LATIN1')</code></td>
<td><code>text_in_utf8</code> represented in Latin-1 encoding (ISO 8859-1)</td>
</tr>
<tr>
<td><code>convert_from(string bytea, src_encoding name)</code></td>
<td>text</td>
<td>Convert string to the data-base encoding. The original encoding is specified by <code>src_encoding</code>. The string</td>
<td><code>convert_from('text_in_utf8', 'UTF8')</code></td>
<td><code>text_in_utf8</code> represented in the current database encoding</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>---------------------------------------------------</td>
<td>--------------------</td>
</tr>
<tr>
<td>convert_to</td>
<td>byte</td>
<td>Convert string to <code>dest_encoding</code>. Must be valid in this encoding.</td>
<td><code>convert_to('some text', 'UTF8')</code></td>
<td>some text represented in the UTF8 encoding</td>
</tr>
<tr>
<td>decode</td>
<td>byte</td>
<td>Decode binary data from textual representation in <code>format</code>. Options 'base64' for <code>format</code> are same as in <code>encode</code>.</td>
<td><code>decode('MTIzAAE=', 'base64')</code></td>
<td><code>\x3132330001</code></td>
</tr>
<tr>
<td>encode</td>
<td>byte</td>
<td>Encode binary data into a textual representation. Supported formats are: base64, hex, escape. escape converts zero bytes and high-bit-set bytes to octal sequences (<code>\nnn</code>) and doubles backslashes.</td>
<td><code>encode('123\000\001', 'base64')</code></td>
<td>MTIzAAE=</td>
</tr>
<tr>
<td>format</td>
<td>text</td>
<td>Format arguments according to a format string. This function is similar to the C function <code>sprintf</code>. See Section 9.4.1.</td>
<td><code>format('Hello %s, %1$s', 'World')</code></td>
<td>Hello World, World</td>
</tr>
<tr>
<td>initcap</td>
<td>text</td>
<td>Convert the first letter of each word to upper case and the rest to lower case. Words are sequences of alphanumeric characters separated by non-alphanumeric characters.</td>
<td><code>initcap('hi THOMAS')</code></td>
<td>Hi Thomas</td>
</tr>
<tr>
<td>left</td>
<td>text</td>
<td>Return first <code>n</code> characters in the string. When <code>n</code> is negative, return all but last `</td>
<td>n</td>
<td>` characters.</td>
</tr>
<tr>
<td>length</td>
<td>int</td>
<td>Number of characters in string.</td>
<td><code>length('jose')</code></td>
<td>4</td>
</tr>
<tr>
<td>length</td>
<td>byte</td>
<td>Number of characters in string in <code>encoding</code>.</td>
<td><code>length('jose', 'UTF8')</code></td>
<td>'UTF8'</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------------</td>
<td>-------------</td>
<td>------------------------------------------------------------------------------</td>
<td>------------------</td>
<td>-----------------</td>
</tr>
<tr>
<td><code>lpad(string text, length [text])</code></td>
<td><code>text</code></td>
<td>Fill up the string to length <code>length</code> by prepending the characters <code>text</code> (a space by default). If the string is already longer than length then it is truncated (on the right).</td>
<td><code>lpad('hi', 5, 'xy')</code></td>
<td><code>xyxhi</code></td>
</tr>
<tr>
<td><code>ltrim(string text [text])</code></td>
<td><code>text</code></td>
<td>Remove the longest string containing only characters from <code>text</code> (a space by default) from the start of <code>string</code>.</td>
<td><code>ltrim('zzzytest', 'xyz')</code></td>
<td><code>test</code></td>
</tr>
<tr>
<td><code>md5(string) text</code></td>
<td><code>text</code></td>
<td>Calculates the MD5 hash of <code>string</code>, returning the result in hexadecimal.</td>
<td><code>md5('abc')</code></td>
<td><code>900150983cd24fb0d6963f7d28e17f72</code></td>
</tr>
<tr>
<td><code>parse_identi-t(string text [text])</code></td>
<td></td>
<td>Split <code>text</code> into an array of identifiers, removing any quoting of individual identifiers. By default, extra characters after the last identifier are considered an error; but if the second parameter is <code>false</code>, then such extra characters are ignored. (This behavior is useful for parsing names for objects like functions.) Note that this function does not truncate over-length identifiers. If you want truncation you can cast the result to <code>name[]</code>.</td>
<td><code>parse_identi-t('SomeSchema'.someTable')</code></td>
<td><code>{SomeSchema,sometable}</code></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>-----------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>------------------------</td>
<td>------------</td>
</tr>
<tr>
<td><code>pg_client_encoding()</code></td>
<td>name</td>
<td>Current client encoding name</td>
<td><code>pg_client_encoding()</code></td>
<td>SQL_ASCII</td>
</tr>
<tr>
<td><code>quote_ident(text)</code></td>
<td>text</td>
<td>Return the given string suitably quoted to be used as an identifier in an SQL statement string. Quotes are added only if necessary (i.e., if the string contains non-identifier characters or would be casefolded). Embedded quotes are properly doubled. See also Example 43.1.</td>
<td><code>quote_ident('Foo bar')</code></td>
<td>&quot;Foo bar&quot;</td>
</tr>
<tr>
<td><code>quote_literal(text)</code></td>
<td>text</td>
<td>Return the given string suitably quoted to be used as a string literal in an SQL statement string; or, if the argument is null, return NULL. Embedded single-quotes and backslashes are properly doubled.</td>
<td><code>quote_literal(E'O\'Reilly')</code></td>
<td>'O''Reilly'</td>
</tr>
<tr>
<td><code>quote_literal(value)</code></td>
<td>text</td>
<td>Coerce the given value to text and then quote it as a literal. Embedded single-quotes and backslashes are properly doubled.</td>
<td><code>quote_literal(42.5)</code></td>
<td>'42.5'</td>
</tr>
<tr>
<td><code>quote_nullable(text)</code></td>
<td>text</td>
<td>Return the given string suitably quoted to be used as a string literal in an SQL statement string; or, if the argument is null, return NULL. Embedded single-quotes and backslashes are properly doubled.</td>
<td><code>quote_nullable(NULL)</code></td>
<td>NULL</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>--------------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------------------------------------------------------------------------</td>
<td>---------------------------------------------------------------------------------------------</td>
<td>-------------------------</td>
</tr>
<tr>
<td>quote_nul-lable(value anyelement)</td>
<td>text</td>
<td>Coerce the given value to text and then quote it as a literal; or, if the argument is NULL, return NULL. Embedded single-quotes and backslashes are properly doubled.</td>
<td>quote_nul-lable(42.5)</td>
<td>'42.5'</td>
</tr>
<tr>
<td>regex-p_match(string text, pattern text [], flags text)</td>
<td>text[]</td>
<td>Return captured substring(s) resulting from the first match of a POSIX regular expression to the string. See Section 9.7.3 for more information.</td>
<td>regex-p_match('foo-bar-bequebaz', '(bar)')</td>
<td>{bar, beque}</td>
</tr>
<tr>
<td>regex-p_matches(string text, pattern text [], flags text)</td>
<td>setof text</td>
<td>Return captured substring(s) resulting from matching a POSIX regular expression to the string. See Section 9.7.3 for more information.</td>
<td>regex-p_matches('foobar-bequebaz', 'ba.', 'g')</td>
<td>{bar}, {baz}</td>
</tr>
<tr>
<td>regexp_replace(string text, pattern text, replacement text [], flags text)</td>
<td>text</td>
<td>Replace substring(s) matching a POSIX regular expression. See Section 9.7.3 for more information.</td>
<td>regexp_replace('Thomas', '.[mN]a.', 'M')</td>
<td>ThM</td>
</tr>
<tr>
<td>regexp_split_to_array(string text, pattern text [])</td>
<td>text[]</td>
<td>Split string using a POSIX regular expression as the delimiter. See Section 9.7.3 for more information.</td>
<td>regexp_split_to_array('hello world', '\s')</td>
<td>{hello, world}</td>
</tr>
<tr>
<td>regexp_split_to_table(string text, pattern text [])</td>
<td>setof text</td>
<td>Split string using a POSIX regular expression as the delimiter. See Section 9.7.3 for more information.</td>
<td>regexp_split_to_table('hello world', '\s')</td>
<td>hello world</td>
</tr>
<tr>
<td>repeat(string text)</td>
<td></td>
<td>Repeat string the specified number of times</td>
<td>repeat('Pg', 4)</td>
<td>PgPgPgPgPg</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------</td>
<td>-------------</td>
<td>---------</td>
<td>----------------</td>
</tr>
<tr>
<td><code>text, number</code></td>
<td><code>int</code></td>
<td>Replace all occurrences in <code>string</code> place with substring from to 'cd', 'XX')</td>
<td><code>re</code>place('abcde-fabcdef', 'cd', 'XX')</td>
<td>abXXefabXXef</td>
</tr>
<tr>
<td><code>text, from text, to text</code></td>
<td><code>text</code></td>
<td>Return reversed string.</td>
<td><code>reverse('abcde')</code></td>
<td>edcba</td>
</tr>
<tr>
<td><code>text, n int)</code></td>
<td><code>text</code></td>
<td>Return last n characters in the string. When n is negative, return all but first</td>
<td><code>right('abcde', 2)</code></td>
<td>de</td>
</tr>
<tr>
<td><code>string text, length int [, fill text])</code></td>
<td><code>text</code></td>
<td>Fill up the string to length by appending the characters fill (a space by default). If the string is already longer than length then it is truncated.</td>
<td><code>rpad('hi', 5, hixyx string to length 'xy')</code></td>
<td>hixyx</td>
</tr>
<tr>
<td><code>string text [, char-acters text])</code></td>
<td><code>text</code></td>
<td>Remove the longest string containing only characters from characters (a space by default) from the end of string</td>
<td><code>rtrim('testxxzx', 'xyz')</code></td>
<td>test</td>
</tr>
<tr>
<td><code>string text, deli-miter text, field int)</code></td>
<td><code>text</code></td>
<td>Split string on delimiter and return the given field (counting from one)</td>
<td><code>split_part('alphabet', '~@~', 2)</code></td>
<td>def</td>
</tr>
<tr>
<td><code>string, substring)</code></td>
<td><code>str</code></td>
<td>Location of specified substring in string, but note the reversed argument order</td>
<td><code>pos('high', 'ig')</code></td>
<td>2</td>
</tr>
<tr>
<td><code>string, from [[, count])</code></td>
<td><code>sub-text</code></td>
<td>Extract substring substring from string (same as substr in string), but note the reversed argument order</td>
<td><code>substr('alphabet', 3, from (string2) from (count))</code></td>
<td></td>
</tr>
<tr>
<td><code>string, prefix)</code></td>
<td><code>bool</code></td>
<td>Returns true if string starts with prefix.</td>
<td><code>starts_with('alphabet', 'alp')</code></td>
<td>t</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_ascii(text)</code></td>
<td><code>text</code></td>
<td>Convert string to ASCII from another encoding (only supports conversion from LATIN1, LATIN2, LATIN9, and WIN1250 encodings)</td>
<td><code>to_ascii('Karel')</code></td>
<td><code>Karel</code></td>
</tr>
<tr>
<td><code>to_hex(number)</code></td>
<td><code>text</code></td>
<td>Convert number to its equivalent hexadecimal representation</td>
<td><code>to_hex(2147483647)</code></td>
<td><code>7fffffff</code></td>
</tr>
<tr>
<td><code>translate(text)</code></td>
<td><code>text</code></td>
<td>Any character in string that matches a character in the from set is replaced by the corresponding character in the to set. If from is longer than to, occurrences of the extra characters in from are removed.</td>
<td><code>translate('12345', '143', 'ax')</code></td>
<td><code>a2x5</code></td>
</tr>
</tbody>
</table>

The `concat`, `concat_ws` and `format` functions are variadic, so it is possible to pass the values to be concatenated or formatted as an array marked with the VARIADIC keyword (see Section 38.5.5). The array’s elements are treated as if they were separate ordinary arguments to the function. If the variadic array argument is NULL, `concat` and `concat_ws` return NULL, but `format` treats a NULL as a zero-element array.

See also the aggregate function `string_agg` in Section 9.20.

### Table 9.10. Built-in Conversions

<table>
<thead>
<tr>
<th>Conversion Name</th>
<th>Source Encoding</th>
<th>Destination Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>ascii_to_mic</td>
<td>SQL_ASCII</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>ascii_to_utf8</td>
<td>SQL_ASCII</td>
<td>UTF8</td>
</tr>
<tr>
<td>big5_to_euc_tw</td>
<td>BIG5</td>
<td>EUC_TW</td>
</tr>
<tr>
<td>big5_to_mic</td>
<td>BIG5</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>big5_to_utf8</td>
<td>BIG5</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_cn_to_mic</td>
<td>EUC_CN</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_cn_to_utf8</td>
<td>EUC_CN</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_jp_to_mic</td>
<td>EUC_JP</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_jp_to_sjis</td>
<td>EUC_JP</td>
<td>SJIS</td>
</tr>
<tr>
<td>euc_jp_to_utf8</td>
<td>EUC_JP</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_kr_to_mic</td>
<td>EUC_KR</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_kr_to_utf8</td>
<td>EUC_KR</td>
<td>UTF8</td>
</tr>
</tbody>
</table>

---

213
<table>
<thead>
<tr>
<th>Conversion Name</th>
<th>Source Encoding</th>
<th>Destination Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>euc_tw_to_big5</td>
<td>EUC_TW</td>
<td>BIG5</td>
</tr>
<tr>
<td>euc_tw_to_mic</td>
<td>EUC_TW</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_tw_to_utf8</td>
<td>EUC_TW</td>
<td>UTF8</td>
</tr>
<tr>
<td>gb18030_to_utf8</td>
<td>GB18030</td>
<td>UTF8</td>
</tr>
<tr>
<td>gbk_to_utf8</td>
<td>GBK</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_10_to_utf8</td>
<td>LATIN6</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_13_to_utf8</td>
<td>LATIN7</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_14_to_utf8</td>
<td>LATIN8</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_15_to_utf8</td>
<td>LATIN9</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_16_to_utf8</td>
<td>LATIN10</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_1_to_mic</td>
<td>LATIN1</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_1_to_utf8</td>
<td>LATIN1</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_2_to_mic</td>
<td>LATIN2</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_2_to_utf8</td>
<td>LATIN2</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_2_to_windows_1250</td>
<td>LATIN2</td>
<td>WIN1250</td>
</tr>
<tr>
<td>iso_8859_3_to_mic</td>
<td>LATIN3</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_3_to_utf8</td>
<td>LATIN3</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_4_to_mic</td>
<td>LATIN4</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_4_to_utf8</td>
<td>LATIN4</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_5_to_koi8_r</td>
<td>ISO_8859_5</td>
<td>KOI8R</td>
</tr>
<tr>
<td>iso_8859_5_to_mic</td>
<td>ISO_8859_5</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_5_to_utf8</td>
<td>ISO_8859_5</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_5_to_windows_1251</td>
<td>ISO_8859_5</td>
<td>WIN1251</td>
</tr>
<tr>
<td>iso_8859_5_to_windows_866</td>
<td>ISO_8859_5</td>
<td>WIN866</td>
</tr>
<tr>
<td>iso_8859_6_to_utf8</td>
<td>ISO_8859_6</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_7_to_utf8</td>
<td>ISO_8859_7</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_8_to_utf8</td>
<td>ISO_8859_8</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_9_to_utf8</td>
<td>LATIN5</td>
<td>UTF8</td>
</tr>
<tr>
<td>johab_to_utf8</td>
<td>JOHAB</td>
<td>UTF8</td>
</tr>
<tr>
<td>koi8_r_to_iso_8859_5</td>
<td>KOI8R</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>koi8_r_to_mic</td>
<td>KOI8R</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>koi8_r_to_utf8</td>
<td>KOI8R</td>
<td>UTF8</td>
</tr>
<tr>
<td>koi8_r_to_windows_1251</td>
<td>KOI8R</td>
<td>WIN1251</td>
</tr>
<tr>
<td>koi8_r_to_windows_866</td>
<td>KOI8R</td>
<td>WIN866</td>
</tr>
<tr>
<td>ko8_u_to_utf8</td>
<td>KOI8U</td>
<td>UTF8</td>
</tr>
<tr>
<td>mic_to_ascii</td>
<td>MULE_INTERNAL</td>
<td>SQL_ASCII</td>
</tr>
<tr>
<td>mic_to_big5</td>
<td>MULE_INTERNAL</td>
<td>BIG5</td>
</tr>
<tr>
<td>mic_to_euc_cn</td>
<td>MULE_INTERNAL</td>
<td>EUC_CN</td>
</tr>
<tr>
<td>Conversion Name</td>
<td>Source Encoding</td>
<td>Destination Encoding</td>
</tr>
<tr>
<td>-------------------------</td>
<td>----------------</td>
<td>----------------------</td>
</tr>
<tr>
<td>mic_to_euc_jp</td>
<td>MULE_INTERNAL</td>
<td>EUC_JP</td>
</tr>
<tr>
<td>mic_to_euc_kr</td>
<td>MULE_INTERNAL</td>
<td>EUC_KR</td>
</tr>
<tr>
<td>mic_to_euc_tw</td>
<td>MULE_INTERNAL</td>
<td>EUC_TW</td>
</tr>
<tr>
<td>mic_to_iso_8859_1</td>
<td>MULE_INTERNAL</td>
<td>LATIN1</td>
</tr>
<tr>
<td>mic_to_iso_8859_2</td>
<td>MULE_INTERNAL</td>
<td>LATIN2</td>
</tr>
<tr>
<td>mic_to_iso_8859_3</td>
<td>MULE_INTERNAL</td>
<td>LATIN3</td>
</tr>
<tr>
<td>mic_to_iso_8859_4</td>
<td>MULE_INTERNAL</td>
<td>LATIN4</td>
</tr>
<tr>
<td>mic_to_iso_8859_5</td>
<td>MULE_INTERNAL</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>mic_to_koi8_r</td>
<td>MULE_INTERNAL</td>
<td>KOI8R</td>
</tr>
<tr>
<td>mic_to_sjis</td>
<td>MULE_INTERNAL</td>
<td>SJIS</td>
</tr>
<tr>
<td>mic_to_windows_1250</td>
<td>MULE_INTERNAL</td>
<td>WIN1250</td>
</tr>
<tr>
<td>mic_to_windows_1251</td>
<td>MULE_INTERNAL</td>
<td>WIN1251</td>
</tr>
<tr>
<td>mic_to_windows_866</td>
<td>MULE_INTERNAL</td>
<td>WIN866</td>
</tr>
<tr>
<td>sjis_to_euc_jp</td>
<td>SJIS</td>
<td>EUC_JP</td>
</tr>
<tr>
<td>sjis_to_mic</td>
<td>SJIS</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>sjis_to_utf8</td>
<td>SJIS</td>
<td>UTF8</td>
</tr>
<tr>
<td>tcvn_to_utf8</td>
<td>WIN1258</td>
<td>UTF8</td>
</tr>
<tr>
<td>uhc_to_utf8</td>
<td>UHC</td>
<td>UTF8</td>
</tr>
<tr>
<td>utf8_to_ascii</td>
<td>UTF8</td>
<td>SQL_ASCII</td>
</tr>
<tr>
<td>utf8_to_big5</td>
<td>UTF8</td>
<td>BIG5</td>
</tr>
<tr>
<td>utf8_to_euc_cn</td>
<td>UTF8</td>
<td>EUC_CN</td>
</tr>
<tr>
<td>utf8_to_euc_jp</td>
<td>UTF8</td>
<td>EUC_JP</td>
</tr>
<tr>
<td>utf8_to_euc_kr</td>
<td>UTF8</td>
<td>EUC_KR</td>
</tr>
<tr>
<td>utf8_to_euc_tw</td>
<td>UTF8</td>
<td>EUC_TW</td>
</tr>
<tr>
<td>utf8_to_gb18030</td>
<td>UTF8</td>
<td>GB18030</td>
</tr>
<tr>
<td>utf8_to_gbk</td>
<td>UTF8</td>
<td>GBK</td>
</tr>
<tr>
<td>utf8_to_iso_8859_1</td>
<td>UTF8</td>
<td>LATIN1</td>
</tr>
<tr>
<td>utf8_to_iso_8859_10</td>
<td>UTF8</td>
<td>LATIN6</td>
</tr>
<tr>
<td>utf8_to_iso_8859_13</td>
<td>UTF8</td>
<td>LATIN7</td>
</tr>
<tr>
<td>utf8_to_iso_8859_14</td>
<td>UTF8</td>
<td>LATIN8</td>
</tr>
<tr>
<td>utf8_to_iso_8859_15</td>
<td>UTF8</td>
<td>LATIN9</td>
</tr>
<tr>
<td>utf8_to_iso_8859_16</td>
<td>UTF8</td>
<td>LATIN10</td>
</tr>
<tr>
<td>utf8_to_iso_8859_2</td>
<td>UTF8</td>
<td>LATIN2</td>
</tr>
<tr>
<td>utf8_to_iso_8859_3</td>
<td>UTF8</td>
<td>LATIN3</td>
</tr>
<tr>
<td>utf8_to_iso_8859_4</td>
<td>UTF8</td>
<td>LATIN4</td>
</tr>
<tr>
<td>utf8_to_iso_8859_5</td>
<td>UTF8</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>utf8_to_iso_8859_6</td>
<td>UTF8</td>
<td>ISO_8859_6</td>
</tr>
<tr>
<td>utf8_to_iso_8859_7</td>
<td>UTF8</td>
<td>ISO_8859_7</td>
</tr>
<tr>
<td>utf8_to_iso_8859_8</td>
<td>UTF8</td>
<td>ISO_8859_8</td>
</tr>
<tr>
<td>utf8_to_iso_8859_9</td>
<td>UTF8</td>
<td>LATIN5</td>
</tr>
<tr>
<td>utf8_to_johab</td>
<td>UTF8</td>
<td>JOHAB</td>
</tr>
<tr>
<td>Conversion Name</td>
<td>Source Encoding</td>
<td>Destination Encoding</td>
</tr>
<tr>
<td>-------------------------</td>
<td>----------------</td>
<td>----------------------</td>
</tr>
<tr>
<td>utf8_to_koi8_r</td>
<td>UTF8</td>
<td>KOI8R</td>
</tr>
<tr>
<td>utf8_to_koi8_u</td>
<td>UTF8</td>
<td>KOI8U</td>
</tr>
<tr>
<td>utf8_to_sjis</td>
<td>UTF8</td>
<td>SJIS</td>
</tr>
<tr>
<td>utf8_to_tcvn</td>
<td>UTF8</td>
<td>WIN1258</td>
</tr>
<tr>
<td>utf8_to_uhc</td>
<td>UTF8</td>
<td>UHC</td>
</tr>
<tr>
<td>utf8_to_windows_1250</td>
<td>UTF8</td>
<td>WIN1250</td>
</tr>
<tr>
<td>utf8_to_windows_1251</td>
<td>UTF8</td>
<td>WIN1251</td>
</tr>
<tr>
<td>utf8_to_windows_1252</td>
<td>UTF8</td>
<td>WIN1252</td>
</tr>
<tr>
<td>utf8_to_windows_1253</td>
<td>UTF8</td>
<td>WIN1253</td>
</tr>
<tr>
<td>utf8_to_windows_1254</td>
<td>UTF8</td>
<td>WIN1254</td>
</tr>
<tr>
<td>utf8_to_windows_1255</td>
<td>UTF8</td>
<td>WIN1255</td>
</tr>
<tr>
<td>utf8_to_windows_1256</td>
<td>UTF8</td>
<td>WIN1256</td>
</tr>
<tr>
<td>utf8_to_windows_1257</td>
<td>UTF8</td>
<td>WIN1257</td>
</tr>
<tr>
<td>utf8_to_windows_866</td>
<td>UTF8</td>
<td>WIN866</td>
</tr>
<tr>
<td>utf8_to_windows_874</td>
<td>UTF8</td>
<td>WIN874</td>
</tr>
<tr>
<td>windows_1250_to_iso_8859_2</td>
<td>WIN1250</td>
<td>LATIN2</td>
</tr>
<tr>
<td>windows_1250_to_mic</td>
<td>WIN1250</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_1250_to_utf8</td>
<td>WIN1250</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_1251_to_iso_8859_5</td>
<td>WIN1251</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>windows_1251_to_koi8_r</td>
<td>WIN1251</td>
<td>KOI8R</td>
</tr>
<tr>
<td>windows_1251_to_mic</td>
<td>WIN1251</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_1251_to_utf8</td>
<td>WIN1251</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_866_to_windows_866</td>
<td>WIN1251</td>
<td>WIN866</td>
</tr>
<tr>
<td>windows_1252_to_utf8</td>
<td>WIN1252</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_1256_to_utf8</td>
<td>WIN1256</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_866_to_iso_8859_5</td>
<td>WIN866</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>windows_866_to_koi8_r</td>
<td>WIN866</td>
<td>KOI8R</td>
</tr>
<tr>
<td>windows_866_to_mic</td>
<td>WIN866</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_866_to_utf8</td>
<td>WIN866</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_866_to_windows_1251</td>
<td>WIN866</td>
<td>WIN</td>
</tr>
<tr>
<td>windows_874_to_utf8</td>
<td>WIN874</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_jis_2004_to_utf8</td>
<td>EUC_JIS_2004</td>
<td>UTF8</td>
</tr>
<tr>
<td>utf8_to_euc_jis_2004</td>
<td>UTF8</td>
<td>EUC_JIS_2004</td>
</tr>
<tr>
<td>shift_jis_2004_to_utf8</td>
<td>SHIFT_JIS_2004</td>
<td>UTF8</td>
</tr>
<tr>
<td>utf8_to_shift_jis_2004</td>
<td>UTF8</td>
<td>SHIFT_JIS_2004</td>
</tr>
</tbody>
</table>
The conversion names follow a standard naming scheme: The official name of the source encoding with all non-alphanumeric characters replaced by underscores, followed by _to_ followed by the similarly processed destination encoding name. Therefore, the names might deviate from the customary encoding names.

### 9.4.1. format

The function `format` produces output formatted according to a format string, in a style similar to the C function `sprintf`.

```c
format(formatstr text [, formatarg "any" [, ...] ])
```

`formatstr` is a format string that specifies how the result should be formatted. Text in the format string is copied directly to the result, except where `format specifiers` are used. `Format specifiers` act as placeholders in the string, defining how subsequent function arguments should be formatted and inserted into the result. Each `formatarg` argument is converted to text according to the usual output rules for its data type, and then formatted and inserted into the result string according to the format specifier(s).

Format specifiers are introduced by a `%` character and have the form

```
%[position][flags][width]type
```

where the component fields are:

- **position** (optional)
  
  A string of the form `n$` where `n` is the index of the argument to print. Index 1 means the first argument after `formatstr`. If the `position` is omitted, the default is to use the next argument in sequence.

- **flags** (optional)
  
  Additional options controlling how the format specifier's output is formatted. Currently the only supported flag is a minus sign (−) which will cause the format specifier's output to be left-justified. This has no effect unless the `width` field is also specified.

- **width** (optional)
  
  Specifies the minimum number of characters to use to display the format specifier's output. The output is padded on the left or right (depending on the − flag) with spaces as needed to fill the width. A too-small width does not cause truncation of the output, but is simply ignored. The width may be specified using any of the following: a positive integer; an asterisk (*) to use the next function argument as the width; or a string of the form `*n$` to use the `n`th function argument as the width.

  If the width comes from a function argument, that argument is consumed before the argument that is used for the format specifier's value. If the width argument is negative, the result is left aligned (as if the − flag had been specified) within a field of length `abs(width)`.

- **type** (required)
  
  The type of format conversion to use to produce the format specifier's output. The following types are supported:
- `s` formats the argument value as a simple string. A null value is treated as an empty string.

- `I` treats the argument value as an SQL identifier, double-quoting it if necessary. It is an error for the value to be null (equivalent to `quote_ident`).

- `L` quotes the argument value as an SQL literal. A null value is displayed as the string `NULL`, without quotes (equivalent to `quote_nullable`).

In addition to the format specifiers described above, the special sequence `%%` may be used to output a literal `%` character.

Here are some examples of the basic format conversions:

```sql
SELECT format('Hello %s', 'World');
Result: Hello World

SELECT format('Testing %s, %s, %s, %s', 'one', 'two', 'three');
Result: Testing one, two, three, %

SELECT format('INSERT INTO %I VALUES(%L)', 'Foo bar', E'O "Reilly');
Result: INSERT INTO "Foo bar" VALUES('O"'Reilly')

SELECT format('INSERT INTO %I VALUES(%L)', 'locations', 'C:\Program Files');
Result: INSERT INTO locations VALUES('C:\Program Files')
```

Here are examples using `width` fields and the `-` flag:

```sql
SELECT format('|%10s|', 'foo');
Result: | foo |

SELECT format('|%-10s|', 'foo');
Result: |foo |

SELECT format('|%*s|', 10, 'foo');
Result: | foo |

SELECT format('|%-*s|', -10, 'foo');
Result: | foo |

SELECT format('|%-*s|', 10, 'foo');
Result: | foo |

SELECT format('|%-*s|', -10, 'foo');
Result: | foo |
```

These examples show use of `position` fields:

```sql
SELECT format('Testing %3$s, %2$s, %1$s', 'one', 'two', 'three');
Result: Testing three, two, one

SELECT format('%*2$s', 'foo', 10, 'bar');
Result: | bar |

SELECT format('%1*$2$s', 'foo', 10, 'bar');
```
Unlike the standard C function `sprintf`, PostgreSQL’s `format` function allows format specifiers with and without position fields to be mixed in the same format string. A format specifier without a position field always uses the next argument after the last argument consumed. In addition, the `format` function does not require all function arguments to be used in the format string. For example:

```sql
SELECT format('Testing %3$s, %2$s, %s', 'one', 'two', 'three');
Result: Testing three, two, three
```

The `%I` and `%L` format specifiers are particularly useful for safely constructing dynamic SQL statements. See Example 43.1.

### 9.5. Binary String Functions and Operators

This section describes functions and operators for examining and manipulating values of type `bytea`.

SQL defines some string functions that use key words, rather than commas, to separate arguments. Details are in Table 9.11. PostgreSQL also provides versions of these functions that use the regular function invocation syntax (see Table 9.12).

#### Note

The sample results shown on this page assume that the server parameter `bytea_output` is set to escape (the traditional PostgreSQL format).

#### Table 9.11. SQL Binary String Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>string</code></td>
<td><code>bytea</code></td>
<td>String concatenation</td>
<td><code>&quot;\Postgres\000\047gres\000&quot;::bytea</code></td>
<td><code>\Postgres\000\047gres\000</code>::bytea</td>
</tr>
<tr>
<td><code>octet_length(string)</code></td>
<td><code>int</code></td>
<td>Number of bytes in binary string</td>
<td><code>octet_length(&quot;jose\000&quot;)::bytea</code></td>
<td>5</td>
</tr>
<tr>
<td><code>overlay(string, place</code></td>
<td><code>bytea</code></td>
<td>Replace substring</td>
<td><code>overlay(&quot;Th\002\003&quot;, \002\003\002::bytea from 2 for 3)</code></td>
<td><code>T\002\003mas</code></td>
</tr>
<tr>
<td><code>position(substring in string)</code></td>
<td><code>int</code></td>
<td>Location of specified substring</td>
<td><code>position('\000om\000', &quot;Th\000omas::bytea&quot;)</code></td>
<td>3</td>
</tr>
<tr>
<td><code>substring(string [from int] [for int])</code></td>
<td><code>bytea</code></td>
<td>Extract substring</td>
<td><code>substring(&quot;Th\000omas::bytea from 2 for 3&quot;)</code></td>
<td><code>h\000o</code></td>
</tr>
<tr>
<td><code>trim([both] bytea bytes from string)</code></td>
<td></td>
<td>Remove the longest string containing only <code>bytea</code></td>
<td><code>trim('\000\001tom\000\001::bytea')</code></td>
<td><code>Tom</code></td>
</tr>
</tbody>
</table>

219
Additional binary string manipulation functions are available and are listed in Table 9.12. Some of them are used internally to implement the SQL-standard string functions listed in Table 9.11.

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>btrim</code></td>
<td><code>bytea</code></td>
<td>Remove the longest string containing only bytes from the start and end of</td>
<td><code>btrim('\000trim\001', 'bytea', \000\001::bytea)</code></td>
<td><code>trim</code></td>
</tr>
<tr>
<td></td>
<td></td>
<td>string</td>
<td></td>
<td><code>123\000456</code></td>
</tr>
<tr>
<td><code>de- code</code></td>
<td><code>bytea</code></td>
<td>Decode binary data from textual representation in 'escape' string. Options</td>
<td><code>de- code('123\000456', 'escape')</code></td>
<td><code>123\000456</code></td>
</tr>
<tr>
<td></td>
<td></td>
<td>for <code>format</code> are same as in <code>encode</code>.</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>encode</code></td>
<td><code>text</code></td>
<td>Encode binary data into a textual representation. Supported formats are:</td>
<td><code>encode('123\000456::bytea', 'escape')</code></td>
<td><code>123\000456</code></td>
</tr>
<tr>
<td></td>
<td></td>
<td>base64, hex, escape. escape converts zero bytes and high-bit-set bytes</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>to octal sequences (<code>\nnn</code>) and doubles backslashes.</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>get_bit</code></td>
<td><code>int</code></td>
<td>Extract bit from string</td>
<td><code>get_bit('Th\000omas', 45)</code></td>
<td><code>1</code></td>
</tr>
<tr>
<td><code>get_byte</code></td>
<td><code>int</code></td>
<td>Extract byte from string</td>
<td><code>get_byte('Th\000omas', 4)</code></td>
<td><code>109</code></td>
</tr>
<tr>
<td><code>length</code></td>
<td><code>int</code></td>
<td>Length of binary string</td>
<td><code>length('jo\000se')</code></td>
<td><code>5</code></td>
</tr>
<tr>
<td><code>md5</code></td>
<td><code>text</code></td>
<td>Calculates the MD5 hash of string, returning the result in hexadecimal</td>
<td><code>md5('Th\0000mas')</code></td>
<td><code>8ab2d3c9689aa18b058c334c82a5b1</code></td>
</tr>
<tr>
<td><code>set_bit</code></td>
<td><code>bytea</code></td>
<td>Set bit in string</td>
<td><code>set_bit('Th\0000mas', 45, 0)</code></td>
<td><code>Th\000omAs</code></td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>set_byte(string, offset, new-value)</td>
<td>bytea</td>
<td>Set byte in string</td>
<td>set_byte('Th\000omas':bytea, 4, 64)</td>
<td></td>
</tr>
<tr>
<td>sha224(bytea)</td>
<td>bytea</td>
<td>SHA-224 hash</td>
<td>sha224('abc') 'x23097d223405d8228642a47da2 55b32aad-bce4b-da0b3f7e36c9da7</td>
<td></td>
</tr>
<tr>
<td>sha256(bytea)</td>
<td>bytea</td>
<td>SHA-256 hash</td>
<td>sha256('abc') 'xba7816bf8f0cfea414140d0b00361a396177a9cb410ff61f20015ad</td>
<td></td>
</tr>
<tr>
<td>sha384(bytea)</td>
<td>bytea</td>
<td>SHA-384 hash</td>
<td>sha384('abc') 'xcb00753f45a25e8b-b5a03d699ac65007 272c32ab0ed-ed1631a8b605a3f-f5bed 8086072bael7c-c2358baec-43c825a7</td>
<td></td>
</tr>
<tr>
<td>sha512(bytea)</td>
<td>bytea</td>
<td>SHA-512 hash</td>
<td>sha512('abc') 'xdaf35a193617abac-c417349ae2041312e6fa4e89a97e2a2a9e1ee64f2192992a274fca836baba3c23a454d4423643ce0e2a9ac94fa0347</td>
<td></td>
</tr>
</tbody>
</table>

get_byte and set_byte number the first byte of a binary string as byte 0. get_bit and set_bit number bits from the right within each byte; for example bit 0 is the least significant bit of the first byte, and bit 15 is the most significant bit of the second byte.

Note that for historic reasons, the function md5 returns a hex-encoded value of type text whereas the SHA-2 functions return type bytea. Use the functions encode and decode to convert between the two, for example encode(sha256('abc'), 'hex') to get a hex-encoded text representation.

See also the aggregate function string_agg in Section 9.20 and the large object functions in Section 35.4.

### 9.6. Bit String Functions and Operators

This section describes functions and operators for examining and manipulating bit strings, that is values of the types bit and bit varying. Aside from the usual comparison operators, the operators shown in Table 9.13 can be used. Bit string operands of & and # must be of equal length. When bit shifting, the original length of the string is preserved, as shown in the examples.

### Table 9.13. Bit String Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>`</td>
<td></td>
<td>`</td>
<td>concatenation</td>
</tr>
<tr>
<td><code>&amp;</code></td>
<td>bitwise AND</td>
<td>`B'10001' &amp; B'01101'</td>
<td><code>00001</code></td>
</tr>
<tr>
<td>`</td>
<td>`</td>
<td>bitwise OR</td>
<td>`B'10001'</td>
</tr>
<tr>
<td>Operator</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>----------------</td>
<td>-------------------</td>
<td>--------</td>
</tr>
<tr>
<td>#</td>
<td>bitwise XOR</td>
<td><code>B'10001'</code></td>
<td><code># 11100</code></td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>B'01101'</code></td>
<td></td>
</tr>
<tr>
<td>~</td>
<td>bitwise NOT</td>
<td>~ <code>B'10001'</code></td>
<td><code>01110</code></td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>bitwise shift left</td>
<td><code>B'10001'</code> &lt;&lt; 3</td>
<td><code>01000</code></td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>bitwise shift right</td>
<td><code>B'10001'</code> &gt;&gt; 2</td>
<td><code>00100</code></td>
</tr>
</tbody>
</table>

The following SQL-standard functions work on bit strings as well as character strings: `length`, `bit_length`, `octet_length`, `position`, `substring`, `overlay`.

The following functions work on bit strings as well as binary strings: `get_bit`, `set_bit`. When working with a bit string, these functions number the first (leftmost) bit of the string as bit 0.

In addition, it is possible to cast integral values to and from type `bit`. Some examples:

- `44::bit(10)`
- `44::bit(3)`
- `cast(-44 as bit(12))`
- `'1110'::bit(4)::integer`

Note that casting to just “bit” means casting to `bit(1)`, and so will deliver only the least significant bit of the integer.

**Note**

Casting an integer to `bit(n)` copies the rightmost n bits. Casting an integer to a bit string width wider than the integer itself will sign-extend on the left.

### 9.7. Pattern Matching

There are three separate approaches to pattern matching provided by PostgreSQL: the traditional SQL `LIKE` operator, the more recent `SIMILAR TO` operator (added in SQL:1999), and POSIX-style regular expressions. Aside from the basic “does this string match this pattern?” operators, functions are available to extract or replace matching substrings and to split a string at matching locations.

**Tip**

If you have pattern matching needs that go beyond this, consider writing a user-defined function in Perl or Tcl.

**Caution**

While most regular-expression searches can be executed very quickly, regular expressions can be contrived that take arbitrary amounts of time and memory to process. Be wary of accepting regular-expression search patterns from hostile sources. If you must do so, it is advisable to impose a statement timeout.

Searches using `SIMILAR TO` patterns have the same security hazards, since `SIMILAR TO` provides many of the same capabilities as POSIX-style regular expressions.
LIKE searches, being much simpler than the other two options, are safer to use with possibly-hostile pattern sources.

9.7.1. LIKE

```sql
string LIKE pattern [ESCAPE escape-character]
string NOT LIKE pattern [ESCAPE escape-character]
```

The LIKE expression returns true if the `string` matches the supplied `pattern`. (As expected, the NOT LIKE expression returns false if LIKE returns true, and vice versa. An equivalent expression is NOT (string LIKE pattern).)

If `pattern` does not contain percent signs or underscores, then the pattern only represents the string itself; in that case LIKE acts like the equals operator. An underscore (_) in `pattern` stands for (matches) any single character; a percent sign (%) matches any sequence of zero or more characters.

Some examples:

```
'abc' LIKE 'abc'    true
'abc' LIKE 'a%'     true
'abc' LIKE '_b_'    true
'abc' LIKE 'c'      false
```

LIKE pattern matching always covers the entire string. Therefore, if it's desired to match a sequence anywhere within a string, the pattern must start and end with a percent sign.

To match a literal underscore or percent sign without matching other characters, the respective character in `pattern` must be preceded by the escape character. The default escape character is the backslash but a different one can be selected by using the `ESCAPE` clause. To match the escape character itself, write two escape characters.

**Note**

If you have `standard_conforming_strings` turned off, any backslashes you write in literal string constants will need to be doubled. See Section 4.1.2.1 for more information.

It's also possible to select no escape character by writing `ESCAPE ''`. This effectively disables the escape mechanism, which makes it impossible to turn off the special meaning of underscore and percent signs in the pattern.

The key word `ILIKE` can be used instead of LIKE to make the match case-insensitive according to the active locale. This is not in the SQL standard but is a PostgreSQL extension.

The operator `~~` is equivalent to `LIKE`, and `~~*` corresponds to `ILIKE`. There are also `!~~` and `!~~*` operators that represent `NOT LIKE` and `NOT ILIKE`, respectively. All of these operators are PostgreSQL-specific.

There is also the prefix operator `^@` and corresponding `starts_with` function which covers cases when only searching by beginning of the string is needed.

9.7.2. SIMILAR TO Regular Expressions

```sql
string SIMILAR TO pattern [ESCAPE escape-character]
```
string NOT SIMILAR TO pattern [ESCAPE escape-character]

The SIMILAR TO operator returns true or false depending on whether its pattern matches the given string. It is similar to LIKE, except that it interprets the pattern using the SQL standard's definition of a regular expression. SQL regular expressions are a curious cross between LIKE notation and common regular expression notation.

Like LIKE, the SIMILAR TO operator succeeds only if its pattern matches the entire string; this is unlike common regular expression behavior where the pattern can match any part of the string. Also like LIKE, SIMILAR TO uses _ and % as wildcard characters denoting any single character and any string, respectively (these are comparable to . and * in POSIX regular expressions).

In addition to these facilities borrowed from LIKE, SIMILAR TO supports these pattern-matching metacharacters borrowed from POSIX regular expressions:

- | denotes alternation (either of two alternatives).
- * denotes repetition of the previous item zero or more times.
- + denotes repetition of the previous item one or more times.
- ? denotes repetition of the previous item zero or one time.
- \{m\} denotes repetition of the previous item exactly \(m\) times.
- \{m, n\} denotes repetition of the previous item \(m\) or more times.
- \{m, n\} denotes repetition of the previous item at least \(m\) and not more than \(n\) times.
- Parentheses () can be used to group items into a single logical item.
- A bracket expression [...] specifies a character class, just as in POSIX regular expressions.

Notice that the period (.) is not a metacharacter for SIMILAR TO.

As with LIKE, a backslash disables the special meaning of any of these metacharacters; or a different escape character can be specified with ESCAPE.

Some examples:

```
'sabc' SIMILAR TO 'abc'     true
'sabc' SIMILAR TO 'a'       false
'sabc' SIMILAR TO '%(b|d)%' true
'sabc' SIMILAR TO '(b|c)%'  false
```

The substring function with three parameters, substring(string from pattern for escape-character), provides extraction of a substring that matches an SQL regular expression pattern. As with SIMILAR TO, the specified pattern must match the entire data string, or else the function fails and returns null. To indicate the part of the pattern that should be returned on success, the pattern must contain two occurrences of the escape character followed by a double quote ("). The text matching the portion of the pattern between these markers is returned.

Some examples, with " delimiting the return string:

```
substring('foobar' from '%"o_b"%' for '#')   oob
substring('foobar' from '#"o_b"%' for '#')   NULL
```

9.7.3. POSIX Regular Expressions

Table 9.14 lists the available operators for pattern matching using POSIX regular expressions.
Table 9.14. Regular Expression Match Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>~</td>
<td>Matches regular expression, case sensitive</td>
<td>'thomas'</td>
<td>.<em>thomas.</em>'</td>
</tr>
<tr>
<td>~*</td>
<td>Matches regular expression, case insensitive</td>
<td>'thomas'</td>
<td>.<em>Thomas.</em>'</td>
</tr>
<tr>
<td>!~</td>
<td>Does not match regular expression, case sensitive</td>
<td>'thomas'</td>
<td>.<em>Thomas.</em>'</td>
</tr>
<tr>
<td>!~*</td>
<td>Does not match regular expression, case insensitive</td>
<td>'thomas'</td>
<td>.<em>vadim.</em>'</td>
</tr>
</tbody>
</table>

POSIX regular expressions provide a more powerful means for pattern matching than the LIKE and SIMILAR TO operators. Many Unix tools such as egrep, sed, or awk use a pattern matching language that is similar to the one described here.

A regular expression is a character sequence that is an abbreviated definition of a set of strings (a regular set). A string is said to match a regular expression if it is a member of the regular set described by the regular expression. As with LIKE, pattern characters match string characters exactly unless they are special characters in the regular expression language — but regular expressions use different special characters than LIKE does. Unlike LIKE patterns, a regular expression is allowed to match anywhere within a string, unless the regular expression is explicitly anchored to the beginning or end of the string.

Some examples:

'abc' ~ 'abc'  true
'abc' ~ '^a'   true
'abc' ~ '(b|d)'  true
'abc' ~ '^(b|c)' false

The POSIX pattern language is described in much greater detail below.

The substring function with two parameters, substring(string from pattern), provides extraction of a substring that matches a POSIX regular expression pattern. It returns null if there is no match, otherwise the portion of the text that matched the pattern. But if the pattern contains any parentheses, the portion of the text that matched the first parenthesized subexpression (the one whose left parenthesis comes first) is returned. You can put parentheses around the whole expression if you want to use parentheses within it without triggering this exception. If you need parentheses in the pattern before the subexpression you want to extract, see the non-capturing parentheses described below.

Some examples:

substring('foobar' from 'o.b')  oob
substring('foobar' from 'o(.b)') o

The regexp_replace function provides substitution of new text for substrings that match POSIX regular expression patterns. It has the syntax regexp_replace(source, pattern, replacement[, flags]). The source string is returned unchanged if there is no match to the pattern. If there is a match, the source string is returned with the replacement string substituted for the matching substring. The replacement string can contain \n, where n is 1 through 9, to indicate that the source substring matching the n'th parenthesized subexpression of the pattern should be inserted, and it can contain \s to indicate that the substring matching the entire pattern should be inserted. Write \\ if you need to put a literal backslash in the replacement text. The flags parameter is an optional text string containing zero or more single-letter flags that change the function’s behavior. Flag i specifies case-insensitive matching, while flag g specifies replacement of each matching substring rather than only the first one. Supported flags (though not g) are described in Table 9.22.
Some examples:

```sql
regexp_replace('foobarbaz', 'b..', 'X')
```
```
fooXbaz
```

```sql
regexp_replace('foobarbaz', 'b..', 'X', 'g')
```
```
fooXX
```

```sql
regexp_replace('foobarbaz', 'b(..)', 'X\1Y', 'g')
```
```
fooXarYXazY
```

The `regexp_match` function returns a text array of captured substring(s) resulting from the first match of a POSIX regular expression pattern to a string. It has the syntax `regexp_match(string, pattern[, flags])`. If there is no match, the result is NULL. If a match is found, and the `pattern` contains no parenthesized subexpressions, then the result is a single-element text array containing the substring matching the whole pattern. If a match is found, and the `pattern` contains parenthesized subexpressions, then the result is a text array whose n'th element is the substring matching the n'th parenthesized subexpression of the `pattern` (not counting “non-capturing” parentheses; see below for details). The `flags` parameter is an optional text string containing zero or more single-letter flags that change the function's behavior. Supported flags are described in Table 9.22.

Some examples:

```sql
SELECT regexp_match('foobarbequebaz', 'bar.*que');
```
```
regexp_match
--------------
{barbeque}
(1 row)
```

```sql
SELECT regexp_match('foobarbequebaz', '(bar)(beque)');
```
```
regexp_match
--------------
{bar,beque}
(1 row)
```

In the common case where you just want the whole matching substring or NULL for no match, write something like

```sql
SELECT (regexp_match('foobarbequebaz', 'bar.*que'))[1];
```
```
regexp_match
--------------
barbeque
(1 row)
```

The `regexp_matches` function returns a set of text arrays of captured substring(s) resulting from matching a POSIX regular expression pattern to a string. It has the same syntax as `regexp_match`. This function returns no rows if there are no rows if there is no match, one row if there is a match and the g flag is not given, or N rows if there are N matches and the g flag is given. Each returned row is a text array containing the whole matched substring or the substrings matching parenthesized subexpressions of the `pattern`, just as described above for `regexp_match`. `regexp_matches` accepts all the flags shown in Table 9.22, plus the g flag which commands it to return all matches, not just the first one.

Some examples:

```sql
SELECT regexp_matches('foo', 'not there');
```
```
regexp_matches
--------------
(0 rows)
```

226
Functions and Operators

SELECT regexp_matches('foobardequebazilbarfbonk', '(b[^b]+) (b[^b]+)', 'g');

regexp_matches
----------------
{bar,beque}
{bazil,barf}
(2 rows)

**Tip**

In most cases `regexp_matches()` should be used with the `g` flag, since if you only want the first match, it's easier and more efficient to use `regexp_match()`. However, `regexp_match()` only exists in PostgreSQL version 10 and up. When working in older versions, a common trick is to place a `regexp_matches()` call in a sub-select, for example:

```
SELECT col1, (SELECT regexp_matches(col2, '(bar) (beque)')) FROM tab;
```

This produces a text array if there's a match, or NULL if not, the same as `regexp_match()` would do. Without the sub-select, this query would produce no output at all for table rows without a match, which is typically not the desired behavior.

The `regexp_split_to_table` function splits a string using a POSIX regular expression pattern as a delimiter. It has the syntax `regexp_split_to_table(string, pattern [, flags])`. If there is no match to the `pattern`, the function returns the `string`. If there is at least one match, for each match it returns the text from the end of the last match (or the beginning of the string) to the beginning of the match. When there are no more matches, it returns the text from the end of the last match to the end of the string. The `flags` parameter is an optional text string containing zero or more single-letter flags that change the function's behavior. `regexp_split_to_table` supports the flags described in Table 9.22.

The `regexp_split_to_array` function behaves the same as `regexp_split_to_table`, except that `regexp_split_to_array` returns its result as an array of text. It has the syntax `regexp_split_to_array(string, pattern [, flags])`. The parameters are the same as for `regexp_split_to_table`.

Some examples:

```
SELECT foo FROM regexp_split_to_table('the quick brown fox jumps over the lazy dog', '\s+') AS foo;
```

```
foo
-----
the
quick
brown
fox
jumps
over
the
lazy
dog
(9 rows)
```
Functions and Operators

SELECT regexp_split_to_array('the quick brown fox jumps over the lazy dog', '\s+');

regexp_split_to_array
-----------------------------------------------
{the, quick, brown, fox, jumps, over, the, lazy, dog}
(1 row)

SELECT foo FROM regexp_split_to_table('the quick brown fox', '\s*')
AS foo;

foo
-----
the quick brown fox
(16 rows)

As the last example demonstrates, the regexp split functions ignore zero-length matches that occur at the start or end of the string or immediately after a previous match. This is contrary to the strict definition of regexp matching that is implemented by regexp_match and regexp_matches, but is usually the most convenient behavior in practice. Other software systems such as Perl use similar definitions.

9.7.3.1. Regular Expression Details

PostgreSQL’s regular expressions are implemented using a software package written by Henry Spencer. Much of the description of regular expressions below is copied verbatim from his manual.

Regular expressions (REs), as defined in POSIX 1003.2, come in two forms: extended REs or EREs (roughly those of egrep), and basic REs or BREs (roughly those of ed). PostgreSQL supports both forms, and also implements some extensions that are not in the POSIX standard, but have become widely used due to their availability in programming languages such as Perl and Tcl. REs using these non-POSIX extensions are called advanced REs or AREs in this documentation. AREs are almost an exact superset of EREs, but BREs have several notational incompatibilities (as well as being much more limited). We first describe the ARE and ERE forms, noting features that apply only to AREs, and then describe how BREs differ.

Note

PostgreSQL always initially presumes that a regular expression follows the ARE rules. However, the more limited ERE or BRE rules can be chosen by prepending an embedded option to the RE pattern, as described in Section 9.7.3.4. This can be useful for compatibility with applications that expect exactly the POSIX 1003.2 rules.
A regular expression is defined as one or more branches, separated by `. It matches anything that matches one of the branches.

A branch is zero or more quantified atoms or constraints, concatenated. It matches a match for the first, followed by a match for the second, etc; an empty branch matches the empty string.

A quantified atom is an atom possibly followed by a single quantifier. Without a quantifier, it matches a match for the atom. With a quantifier, it can match some number of matches of the atom. An atom can be any of the possibilities shown in Table 9.15. The possible quantifiers and their meanings are shown in Table 9.16.

A constraint matches an empty string, but matches only when specific conditions are met. A constraint can be used where an atom could be used, except it cannot be followed by a quantifier. The simple constraints are shown in Table 9.17; some more constraints are described later.

### Table 9.15. Regular Expression Atoms

<table>
<thead>
<tr>
<th>Atom</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>(re)</code></td>
<td>(where <code>re</code> is any regular expression) matches a match for <code>re</code>, with the match noted for possible reporting</td>
</tr>
<tr>
<td><code>(?:re)</code></td>
<td>as above, but the match is not noted for reporting (a “non-capturing” set of parentheses) (AREs only)</td>
</tr>
<tr>
<td><code>.</code></td>
<td>matches any single character</td>
</tr>
<tr>
<td><code>[chars]</code></td>
<td>a bracket expression, matching any one of the <code>chars</code> (see Section 9.7.3.2 for more detail)</td>
</tr>
<tr>
<td><code>\k</code></td>
<td>(where <code>k</code> is a non-alphanumeric character) matches that character taken as an ordinary character, e.g., <code>\</code> matches a backslash character</td>
</tr>
<tr>
<td><code>\c</code></td>
<td>where <code>c</code> is alphanumeric (possibly followed by other characters) is an escape, see Section 9.7.3.3 (AREs only; in EREs and BREs, this matches <code>c</code>)</td>
</tr>
<tr>
<td><code>{</code></td>
<td>when followed by a character other than a digit, matches the left-brace character <code>{</code>; when followed by a digit, it is the beginning of a bound (see below)</td>
</tr>
<tr>
<td><code>x</code></td>
<td>where <code>x</code> is a single character with no other significance, matches that character</td>
</tr>
</tbody>
</table>

An RE cannot end with a backslash (`\`).

### Note

If you have `standard_conforming_strings` turned off, any backslashes you write in literal string constants will need to be doubled. See Section 4.1.2.1 for more information.

### Table 9.16. Regular Expression Quantifiers

<table>
<thead>
<tr>
<th>Quantifier</th>
<th>Matches</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>*</code></td>
<td>a sequence of 0 or more matches of the atom</td>
</tr>
<tr>
<td><code>+</code></td>
<td>a sequence of 1 or more matches of the atom</td>
</tr>
</tbody>
</table>
Functions and Operators

<table>
<thead>
<tr>
<th>Quantifier</th>
<th>Matches</th>
</tr>
</thead>
<tbody>
<tr>
<td>?</td>
<td>a sequence of 0 or 1 matches of the atom</td>
</tr>
<tr>
<td>{ m }</td>
<td>a sequence of exactly m matches of the atom</td>
</tr>
<tr>
<td>{ m, }</td>
<td>a sequence of m or more matches of the atom</td>
</tr>
<tr>
<td>{ m, n }</td>
<td>a sequence of m through n (inclusive) matches of the atom; n cannot exceed m</td>
</tr>
<tr>
<td>*?</td>
<td>non-greedy version of *</td>
</tr>
<tr>
<td>+?</td>
<td>non-greedy version of +</td>
</tr>
<tr>
<td>??</td>
<td>non-greedy version of ?</td>
</tr>
<tr>
<td>{ m }?</td>
<td>non-greedy version of { m }</td>
</tr>
<tr>
<td>{ m, }?</td>
<td>non-greedy version of { m, }</td>
</tr>
<tr>
<td>{ m, n }?</td>
<td>non-greedy version of { m, n }</td>
</tr>
</tbody>
</table>

The forms using { ... } are known as *bounds*. The numbers m and n within a bound are unsigned decimal integers with permissible values from 0 to 255 inclusive.

Non-greedy quantifiers (available in AREs only) match the same possibilities as their corresponding normal (greedy) counterparts, but prefer the smallest number rather than the largest number of matches. See Section 9.7.3.5 for more detail.

**Note**

A quantifier cannot immediately follow another quantifier, e.g., ** is invalid. A quantifier cannot begin an expression or subexpression or follow ^ or |.

<table>
<thead>
<tr>
<th>Constraint</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>^</td>
<td>matches at the beginning of the string</td>
</tr>
<tr>
<td>$</td>
<td>matches at the end of the string</td>
</tr>
<tr>
<td>(?)=xe</td>
<td><em>positive lookahead</em> matches at any point where a substring matching xe begins (AREs only)</td>
</tr>
<tr>
<td>(?!xe)</td>
<td><em>negative lookahead</em> matches at any point where no substring matching xe begins (AREs only)</td>
</tr>
<tr>
<td>(?&gt;=xe)</td>
<td><em>positive lookbehind</em> matches at any point where a substring matching xe ends (AREs only)</td>
</tr>
<tr>
<td>(&lt;?!xe)</td>
<td><em>negative lookbehind</em> matches at any point where no substring matching xe ends (AREs only)</td>
</tr>
</tbody>
</table>

Lookahead and lookbehind constraints cannot contain back references (see Section 9.7.3.3), and all parentheses within them are considered non-capturing.

### 9.7.3.2. Bracket Expressions

A *bracket expression* is a list of characters enclosed in [ ]. It normally matches any single character from the list (but see below). If the list begins with ^, it matches any single character not from the rest of the list. If two characters in the list are separated by –, this is shorthand for the full range of characters between those two (inclusive) in the collating sequence, e.g., [0–9] in ASCII matches any decimal digit. It is illegal for two ranges to share an endpoint, e.g., a–c–e. Ranges are very collating-sequence-dependent, so portable programs should avoid relying on them.
To include a literal } in the list, make it the first character (after ^, if that is used). To include a literal -, make it the first or last character, or the second endpoint of a range. To use a literal - as the first endpoint of a range, enclose it in [. and .] to make it a collating element (see below).

With the exception of these characters, some combinations using [ (see next paragraphs), and escapes (AREs only), all other special characters lose their special significance within a bracket expression. In particular, \ is not special when following ERE or BRE rules, though it is special (as introducing an escape) in AREs.

Within a bracket expression, a collating element (a character, a multiple-character sequence that collates as if it were a single character, or a collating-sequence name for either) enclosed in [ . and .] stands for the sequence of characters of that collating element. The sequence is treated as a single element of the bracket expression's list. This allows a bracket expression containing a multiple-character collating element to match more than one character, e.g., if the collating sequence includes a ch collating element, then the RE [[.ch.]]*c matches the first five characters of chchcc.

Within a bracket expression, a collating element enclosed in [ = and =] is an equivalence class, standing for the sequences of characters of all collating elements equivalent to that one, including itself. (If there are no other equivalent collating elements, the treatment is as if the enclosing delimiters were [ . and .].) For example, if o and ^ are the members of an equivalence class, then [[=o=]], [[=^=]], and [o^] are all synonymous. An equivalence class cannot be an endpoint of a range.

Within a bracket expression, the name of a character class enclosed in [ : and :] stands for the list of all characters belonging to that class. Standard character class names are: alnum, alpha, blank, cntrl, digit, graph, lower, print, punct, space, upper, xdigit. These stand for the character classes defined in ctype. A locale can provide others. A character class cannot be used as an endpoint of a range.

There are two special cases of bracket expressions: the bracket expressions [[:<:]] and [[:<:>]] are constraints, matching empty strings at the beginning and end of a word respectively. A word is defined as a sequence of word characters that is neither preceded nor followed by word characters. A word character is an alnum character (as defined by ctype) or an underscore. This is an extension, compatible with but not specified by POSIX 1003.2, and should be used with caution in software intended to be portable to other systems. The constraint escapes described below are usually preferable; they are no more standard, but are easier to type.

### 9.7.3.3. Regular Expression Escapes

Escapes are special sequences beginning with \ followed by an alphanumeric character. Escapes come in several varieties: character entry, class shorthands, constraint escapes, and back references. A \ followed by an alphanumeric character but not constituting a valid escape is illegal in AREs. In EREs, there are no escapes: outside a bracket expression, a \ followed by an alphanumeric character merely stands for that character as an ordinary character, and inside a bracket expression, \ is an ordinary character. (The latter is the one actual incompatibility between EREs and AREs.)

**Character-entry escapes** exist to make it easier to specify non-printing and other inconvenient characters in REs. They are shown in Table 9.18.

**Class-shorthand escapes** provide shorthands for certain commonly-used character classes. They are shown in Table 9.19.

A **constraint escape** is a constraint, matching the empty string if specific conditions are met, written as an escape. They are shown in Table 9.20.
A back reference (\n) matches the same string matched by the previous parenthesized subexpression specified by the number \n (see Table 9.21). For example, (\[bc\])\1 matches \bb or \cc but not \bc or \cb. The subexpression must entirely precede the back reference in the RE. Subexpressions are numbered in the order of their leading parentheses. Non-capturing parentheses do not define subexpressions.

Table 9.18. Regular Expression Character-entry Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\a</td>
<td>alert (bell) character, as in C</td>
</tr>
<tr>
<td>\b</td>
<td>backspace, as in C</td>
</tr>
<tr>
<td>\B</td>
<td>synonym for backslash () to help reduce the need for backslash doubling</td>
</tr>
<tr>
<td>\cX</td>
<td>(where \X is any character) the character whose low-order 5 bits are the same as those of \X, and whose other bits are all zero</td>
</tr>
<tr>
<td>\e</td>
<td>the character whose collating-sequence name is ESC, or failing that, the character with octal value 033</td>
</tr>
<tr>
<td>\f</td>
<td>form feed, as in C</td>
</tr>
<tr>
<td>\n</td>
<td>newline, as in C</td>
</tr>
<tr>
<td>\r</td>
<td>carriage return, as in C</td>
</tr>
<tr>
<td>\t</td>
<td>horizontal tab, as in C</td>
</tr>
<tr>
<td>\uwxyz</td>
<td>(where wxyz is exactly four hexadecimal digits) the character whose hexadecimal value is 0xwxyz</td>
</tr>
<tr>
<td>\Ustuvwxyz</td>
<td>(where stuvwxyz is exactly eight hexadecimal digits) the character whose hexadecimal value is 0xstuvwxyz</td>
</tr>
<tr>
<td>\v</td>
<td>vertical tab, as in C</td>
</tr>
<tr>
<td>\xhhh</td>
<td>(where hhh is any sequence of hexadecimal digits) the character whose hexadecimal value is 0xhhhh (a single character no matter how many hexadecimal digits are used)</td>
</tr>
<tr>
<td>\0</td>
<td>the character whose value is 0 (the null byte)</td>
</tr>
<tr>
<td>\xy</td>
<td>(where xy is exactly two octal digits, and is not a back reference) the character whose octal value is 0xy</td>
</tr>
<tr>
<td>\xyz</td>
<td>(where xyz is exactly three octal digits, and is not a back reference) the character whose octal value is 0xyz</td>
</tr>
</tbody>
</table>

Hexadecimal digits are 0-9, a-f, and A-F. Octal digits are 0-7.

Numeric character-entry escapes specifying values outside the ASCII range (0-127) have meanings dependent on the database encoding. When the encoding is UTF-8, escape values are equivalent to Unicode code points, for example \u1234 means the character U+1234. For other multibyte encodings, character-entry escapes usually just specify the concatenation of the byte values for the character. If the escape value does not correspond to any legal character in the database encoding, no error will be raised, but it will never match any data.

The character-entry escapes are always taken as ordinary characters. For example, \135 is ] in ASCII, but \135 does not terminate a bracket expression.
Table 9.19. Regular Expression Class-shorthand Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\d</td>
<td>[:digit:]</td>
</tr>
<tr>
<td>\s</td>
<td>[:space:]</td>
</tr>
<tr>
<td>\w</td>
<td>[:alnum:]_</td>
</tr>
<tr>
<td>\D</td>
<td>^[:digit:]</td>
</tr>
<tr>
<td>\S</td>
<td>^[:space:]</td>
</tr>
<tr>
<td>\W</td>
<td>^[:alnum:]_</td>
</tr>
</tbody>
</table>

Within bracket expressions, \d, \s, and \w lose their outer brackets, and \D, \S, and \W are illegal. (So, for example, [a-c\d] is equivalent to [a-c[:digit:]]. Also, [a-c\D], which is equivalent to [a-c^[:digit:]], is illegal.)

Table 9.20. Regular Expression Constraint Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\A</td>
<td>matches only at the beginning of the string (see Section 9.7.3.5 for how this differs from ^)</td>
</tr>
<tr>
<td>\m</td>
<td>matches only at the beginning of a word</td>
</tr>
<tr>
<td>\M</td>
<td>matches only at the end of a word</td>
</tr>
<tr>
<td>\y</td>
<td>matches only at the beginning or end of a word</td>
</tr>
<tr>
<td>\Y</td>
<td>matches only at a point that is not the beginning or end of a word</td>
</tr>
<tr>
<td>\Z</td>
<td>matches only at the end of the string (see Section 9.7.3.5 for how this differs from $)</td>
</tr>
</tbody>
</table>

A word is defined as in the specification of [:<:] and [:>:] above. Constraint escapes are illegal within bracket expressions.

Table 9.21. Regular Expression Back References

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\m</td>
<td>(where (m) is a nonzero digit) a back reference to the (m)'th subexpression</td>
</tr>
<tr>
<td>\mnn</td>
<td>(where (m) is a nonzero digit, and (nn) is some more digits, and the decimal value (mnn) is not greater than the number of closing capturing parentheses seen so far) a back reference to the (mnn)'th subexpression</td>
</tr>
</tbody>
</table>

Note

There is an inherent ambiguity between octal character-entry escapes and back references, which is resolved by the following heuristics, as hinted at above. A leading zero always indicates an octal escape. A single non-zero digit, not followed by another digit, is always taken as a back reference. A multi-digit sequence not starting with a zero is taken as a back reference if it comes after a suitable subexpression (i.e., the number is in the legal range for a back reference), and otherwise is taken as octal.
9.7.3.4. Regular Expression Metasyntax

In addition to the main syntax described above, there are some special forms and miscellaneous syntactic facilities available.

An RE can begin with one of two special director prefixes. If an RE begins with ***, the rest of the RE is taken as an ARE. (This normally has no effect in PostgreSQL, since REs are assumed to be AREs; but it does have an effect if ERE or BRE mode had been specified by the flags parameter to a regex function.) If an RE begins with ***, the rest of the RE is taken to be a literal string, with all characters considered ordinary characters.

An ARE can begin with embedded options: a sequence (?xyz) (where xyz is one or more alphabetic characters) specifies options affecting the rest of the RE. These options override any previously determined options — in particular, they can override the case-sensitivity behavior implied by a regex operator, or the flags parameter to a regex function. The available option letters are shown in Table 9.22. Note that these same option letters are used in the flags parameters of regex functions.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>b</td>
<td>rest of RE is a BRE</td>
</tr>
<tr>
<td>c</td>
<td>case-sensitive matching (overrides operator type)</td>
</tr>
<tr>
<td>e</td>
<td>rest of RE is an ERE</td>
</tr>
<tr>
<td>i</td>
<td>case-insensitive matching (see Section 9.7.3.5) (overrides operator type)</td>
</tr>
<tr>
<td>m</td>
<td>historical synonym for n</td>
</tr>
<tr>
<td>n</td>
<td>newline-sensitive matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>p</td>
<td>partial newline-sensitive matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>q</td>
<td>rest of RE is a literal (&quot;quoted&quot;) string, all ordinary characters</td>
</tr>
<tr>
<td>s</td>
<td>non-newline-sensitive matching (default)</td>
</tr>
<tr>
<td>t</td>
<td>tight syntax (default; see below)</td>
</tr>
<tr>
<td>w</td>
<td>inverse partial newline-sensitive (&quot;weird&quot;) matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>x</td>
<td>expanded syntax (see below)</td>
</tr>
</tbody>
</table>

Embedded options take effect at the ) terminating the sequence. They can appear only at the start of an ARE (after the **: director if any).

In addition to the usual (tight) RE syntax, in which all characters are significant, there is an expanded syntax, available by specifying the embedded x option. In the expanded syntax, white-space characters in the RE are ignored, as are all characters between a # and the following newline (or the end of the RE). This permits paragraphing and commenting a complex RE. There are three exceptions to that basic rule:

• a white-space character or # preceded by \ is retained

• white space or # within a bracket expression is retained

• white space and comments cannot appear within multi-character symbols, such as (? :

For this purpose, white-space characters are blank, tab, newline, and any character that belongs to the space character class.
Finally, in an ARE, outside bracket expressions, the sequence (?#ttt) (where ttt is any text not containing a ) is a comment, completely ignored. Again, this is not allowed between the characters of multi-character symbols, like (?:. Such comments are more a historical artifact than a useful facility, and their use is deprecated; use the expanded syntax instead.

None of these metasyntax extensions is available if an initial ***= director has specified that the user's input be treated as a literal string rather than as an RE.

9.7.3.5. Regular Expression Matching Rules

In the event that an RE could match more than one substring of a given string, the RE matches the one starting earliest in the string. If the RE could match more than one substring starting at that point, either the longest possible match or the shortest possible match will be taken, depending on whether the RE is greedy or non-greedy.

Whether an RE is greedy or not is determined by the following rules:

- Most atoms, and all constraints, have no greediness attribute (because they cannot match variable amounts of text anyway).
- Adding parentheses around an RE does not change its greediness.
- A quantified atom with a fixed-repetition quantifier (\{m\} or \{m\}?\) has the same greediness (possibly none) as the atom itself.
- A quantified atom with other normal quantifiers (including \{m, n\} with m equal to n) is greedy (prefers longest match).
- A quantified atom with a non-greedy quantifier (including \{m, n\}?\ with m equal to n) is non-greedy (prefers shortest match).
- A branch — that is, an RE that has no top-level \| operator — has the same greediness as the first quantified atom in it that has a greediness attribute.
- An RE consisting of two or more branches connected by the \| operator is always greedy.

The above rules associate greediness attributes not only with individual quantified atoms, but with branches and entire REs that contain quantified atoms. What that means is that the matching is done in such a way that the branch, or whole RE, matches the longest or shortest possible substring as a whole. Once the length of the entire match is determined, the part of it that matches any particular subexpression is determined on the basis of the greediness attribute of that subexpression, with subexpressions starting earlier in the RE taking priority over ones starting later.

An example of what this means:

```
SELECT SUBSTRING('XY1234Z', 'Y*(\[0-9\]{1,3})');
Result: 123
SELECT SUBSTRING('XY1234Z', 'Y*?(\[0-9\]{1,3})');
Result: 1
```

In the first case, the RE as a whole is greedy because \* is greedy. It can match beginning at the Y, and it matches the longest possible string starting there, i.e., \Y123. The output is the parenthesized part of that, or 123. In the second case, the RE as a whole is non-greedy because \*? is non-greedy. It can match beginning at the Y, and it matches the shortest possible string starting there, i.e., Y1. The subexpression \{0–9\}{1,3} is greedy but it cannot change the decision as to the overall match length; so it is forced to match just 1.

In short, when an RE contains both greedy and non-greedy subexpressions, the total match length is either as long as possible or as short as possible, according to the attribute assigned to the whole RE. The attributes assigned to the subexpressions only affect how much of that match they are allowed to "eat" relative to each other.
The quantifiers \( \{1,1\} \) and \( \{1,1\} \) can be used to force greediness or non-greediness, respectively, on a subexpression or a whole RE. This is useful when you need the whole RE to have a greediness attribute different from what’s deduced from its elements. As an example, suppose that we are trying to separate a string containing some digits into the digits and the parts before and after them. We might try to do that like this:

```sql
SELECT regexp_match('abc01234xyz', '(.*)(\d+)(.*)');
```

Result: \{abc0123,4,xyz\}

That didn’t work: the first .* is greedy so it “eats” as much as it can, leaving the \d+ to match at the last possible place, the last digit. We might try to fix that by making it non-greedy:

```sql
SELECT regexp_match('abc01234xyz', '(.*?)(\d+)(.*)');
```

Result: \{abc,0,”\}

That didn’t work either, because now the RE as a whole is non-greedy and so it ends the overall match as soon as possible. We can get what we want by forcing the RE as a whole to be greedy:

```sql
SELECT regexp_match('abc01234xyz', '(?:(.*?)(\d+)(.*)){1,1}');
```

Result: \{abc,01234,xyz\}

Controlling the RE’s overall greediness separately from its components’ greediness allows great flexibility in handling variable-length patterns.

When deciding what is a longer or shorter match, match lengths are measured in characters, not collating elements. An empty string is considered longer than no match at all. For example: \( bb^* \) matches the three middle characters of \abc\b; \( \text{week|wee} \text{night|knights} \) matches all ten characters of \text{weeknights}; when \( .* \). \* is matched against \abc\ the parenthesized subexpression matches all three characters; and when \( (a^*) \) \* is matched against \bc\ both the whole RE and the parenthesized subexpression match an empty string.

If case-independent matching is specified, the effect is much as if all case distinctions had vanished from the alphabet. When an alphabetic that exists in multiple cases appears as an ordinary character outside a bracket expression, it is effectively transformed into a bracket expression containing both cases, e.g., \x\ becomes \[xX\]. When it appears inside a bracket expression, all case counterparts of it are added to the bracket expression, e.g., \[x\] becomes \[xX\] and \[^x\] becomes \[^xX\].

If newline-sensitive matching is specified, \. and bracket expressions using \^ will never match the newline character (so that matches will never cross newlines unless the RE explicitly arranges it) and \^ and \$ will match the empty string after and before a newline respectively, in addition to matching at beginning and end of string respectively. But the ARE escapes \A\ and \Z\ continue to match beginning or end of string only.

If partial newline-sensitive matching is specified, this affects \. and bracket expressions as with newline-sensitive matching, but not \^ and \$.

If inverse partial newline-sensitive matching is specified, this affects \^ and \$ as with newline-sensitive matching, but not \. and bracket expressions. This isn’t very useful but is provided for symmetry.

### 9.7.3.6. Limits and Compatibility

No particular limit is imposed on the length of REs in this implementation. However, programs intended to be highly portable should not employ REs longer than 256 bytes, as a POSIX-compliant implementation can refuse to accept such REs.

The only feature of AREs that is actually incompatible with POSIX EREs is that \ does not lose its special significance inside bracket expressions. All other ARE features use syntax which is illegal or
has undefined or unspecified effects in POSIX EREs; the *** syntax of directors likewise is outside the POSIX syntax for both BREs and EREs.

Many of the ARE extensions are borrowed from Perl, but some have been changed to clean them up, and a few Perl extensions are not present. Incompatibilities of note include \b, \B, the lack of special treatment for a trailing newline, the addition of complemented bracket expressions to the things affected by newline-sensitive matching, the restrictions on parentheses and back references in lookahead/lookbehind constraints, and the longest/shortest-match (rather than first-match) matching semantics.

Two significant incompatibilities exist between AREs and the ERE syntax recognized by pre-7.4 releases of PostgreSQL:

- In AREs, \ followed by an alphanumerical character is either an escape or an error, while in previous releases, it was just another way of writing the alphanumerical. This should not be much of a problem because there was no reason to write such a sequence in earlier releases.

- In AREs, \ remains a special character within [], so a literal \ within a bracket expression must be written \\.

### 9.7.3.7. Basic Regular Expressions

BREs differ from EREs in several respects. In BREs, |, +, and ? are ordinary characters and there is no equivalent for their functionality. The delimiters for bounds are \{ and \}, with ( and ) by themselves ordinary characters. The parentheses for nested subexpressions are \( and \), with ( and ) by themselves ordinary characters. ^ is an ordinary character except at the beginning of the RE or the beginning of a parenthesized subexpression, 5 is an ordinary character except at the end of the RE or the end of a parenthesized subexpression, and * is an ordinary character if it appears at the beginning of the RE or the beginning of a parenthesized subexpression (after a possible leading ^). Finally, single-digit back references are available, and \< and \> are synonyms for [[::]] and [[::]] respectively; no other escapes are available in BREs.

### 9.8. Data Type Formatting Functions

The PostgreSQL formatting functions provide a powerful set of tools for converting various data types (date/time, integer, floating point, numeric) to formatted strings and for converting from formatted strings to specific data types. Table 9.23 lists them. These functions all follow a common calling convention: the first argument is the value to be formatted and the second argument is a template that defines the output or input format.

#### Table 9.23. Formatting Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_char(time-stamp, text)</code></td>
<td>text</td>
<td>convert time stamp to string</td>
<td><code>to_char(current_timestamp, 'HH12:MI:SS')</code></td>
</tr>
<tr>
<td><code>to_char(interval, text)</code></td>
<td>text</td>
<td>convert interval to string</td>
<td><code>to_char(interval '15h 2m 12s', 'HH24:MI:SS')</code></td>
</tr>
<tr>
<td><code>to_char(int, text)</code></td>
<td>text</td>
<td>convert integer to string</td>
<td><code>to_char(125, '999')</code></td>
</tr>
<tr>
<td><code>to_char(double precision, text)</code></td>
<td>text</td>
<td>convert real/double precision to string</td>
<td><code>to_char(125.8::real, '999D99S')</code></td>
</tr>
<tr>
<td><code>to_char(numeric, text)</code></td>
<td>text</td>
<td>convert numeric to string</td>
<td><code>to_char(-125.8, '999D99S')</code></td>
</tr>
<tr>
<td><code>to_date(text, text)</code></td>
<td>date</td>
<td>convert string to date</td>
<td><code>to_date('D8000', 'DD Mon YYYY')</code></td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_number(text, text)</code></td>
<td>numeric</td>
<td>convert string to numeric</td>
<td><code>to_number('12,454.8-', '99G999D9S')</code></td>
</tr>
<tr>
<td><code>to_timestamp(text, text)</code></td>
<td>timestamp with time zone</td>
<td>convert string to timestamp</td>
<td><code>to_timestamp('05 Dec 2000', 'DD Mon YYYY')</code></td>
</tr>
</tbody>
</table>

**Note**

There is also a single-argument `to_timestamp` function; see Table 9.30.

**Tip**

`to_timestamp` and `to_date` exist to handle input formats that cannot be converted by simple casting. For most standard date/time formats, simply casting the source string to the required data type works, and is much easier. Similarly, `to_number` is unnecessary for standard numeric representations.

In a `to_char` output template string, there are certain patterns that are recognized and replaced with appropriately-formatted data based on the given value. Any text that is not a template pattern is simply copied verbatim. Similarly, in an input template string (for the other functions), template patterns identify the values to be supplied by the input data string. If there are characters in the template string that are not template patterns, the corresponding characters in the input data string are simply skipped over (whether or not they are equal to the template string characters).

Table 9.24 shows the template patterns available for formatting date and time values.

**Table 9.24. Template Patterns for Date/Time Formatting**

<table>
<thead>
<tr>
<th>Pattern</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>HH</td>
<td>hour of day (01-12)</td>
</tr>
<tr>
<td>HH12</td>
<td>hour of day (01-12)</td>
</tr>
<tr>
<td>HH24</td>
<td>hour of day (00-23)</td>
</tr>
<tr>
<td>MI</td>
<td>minute (00-59)</td>
</tr>
<tr>
<td>SS</td>
<td>second (00-59)</td>
</tr>
<tr>
<td>MS</td>
<td>millisecond (000-999)</td>
</tr>
<tr>
<td>US</td>
<td>microsecond (000000-999999)</td>
</tr>
<tr>
<td>SSSS</td>
<td>seconds past midnight (0-86399)</td>
</tr>
<tr>
<td>AM, am, PM or pm</td>
<td>meridiem indicator (without periods)</td>
</tr>
<tr>
<td>A.M., a.m., P.M. or p.m.</td>
<td>meridiem indicator (with periods)</td>
</tr>
<tr>
<td>Y, YYYY</td>
<td>year (4 or more digits) with comma</td>
</tr>
<tr>
<td>YYYY</td>
<td>year (4 or more digits)</td>
</tr>
<tr>
<td>YYY</td>
<td>last 3 digits of year</td>
</tr>
<tr>
<td>YY</td>
<td>last 2 digits of year</td>
</tr>
<tr>
<td>Y</td>
<td>last digit of year</td>
</tr>
<tr>
<td>IYYYY</td>
<td>ISO 8601 week-numbering year (4 or more digits)</td>
</tr>
<tr>
<td>IYY</td>
<td>last 3 digits of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>Pattern</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>IY</td>
<td>last 2 digits of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>I</td>
<td>last digit of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>BC, bc, AD or ad</td>
<td>era indicator (without periods)</td>
</tr>
<tr>
<td>B.C., b.c., A.D. or a.d.</td>
<td>era indicator (with periods)</td>
</tr>
<tr>
<td>MONTH</td>
<td>full upper case month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>MONTH</td>
<td>full capitalized month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>month</td>
<td>full lower case month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>MON</td>
<td>abbreviated upper case month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>Mon</td>
<td>abbreviated capitalized month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>mon</td>
<td>abbreviated lower case month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>MM</td>
<td>month number (01-12)</td>
</tr>
<tr>
<td>DAY</td>
<td>full upper case day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>Day</td>
<td>full capitalized day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>day</td>
<td>full lower case day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>DY</td>
<td>abbreviated upper case day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>Dy</td>
<td>abbreviated capitalized day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>dy</td>
<td>abbreviated lower case day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>DDD</td>
<td>day of year (001-366)</td>
</tr>
<tr>
<td>IDDD</td>
<td>day of ISO 8601 week-numbering year (001-371; day 1 of the year is Monday of the first ISO week)</td>
</tr>
<tr>
<td>DD</td>
<td>day of month (01-31)</td>
</tr>
<tr>
<td>D</td>
<td>day of the week, Sunday (1) to Saturday (7)</td>
</tr>
<tr>
<td>ID</td>
<td>ISO 8601 day of the week, Monday (1) to Sunday (7)</td>
</tr>
<tr>
<td>W</td>
<td>week of month (1-5) (the first week starts on the first day of the month)</td>
</tr>
<tr>
<td>WW</td>
<td>week number of year (1-53) (the first week starts on the first day of the year)</td>
</tr>
<tr>
<td>IW</td>
<td>week number of ISO 8601 week-numbering year (01-53; the first Thursday of the year in week 1)</td>
</tr>
<tr>
<td>CC</td>
<td>century (2 digits) (the twenty-first century starts on 2001-01-01)</td>
</tr>
<tr>
<td>J</td>
<td>Julian Day (integer days since November 24, 4714 BC at midnight UTC)</td>
</tr>
<tr>
<td>Pattern</td>
<td>Description</td>
</tr>
<tr>
<td>---------</td>
<td>-------------</td>
</tr>
<tr>
<td>Q</td>
<td>quarter</td>
</tr>
<tr>
<td>RM</td>
<td>month in upper case Roman numerals (I-XII; I=January)</td>
</tr>
<tr>
<td>rm</td>
<td>month in lower case Roman numerals (i-xii; i=January)</td>
</tr>
<tr>
<td>TZ</td>
<td>upper case time-zone abbreviation (only supported in to_char)</td>
</tr>
<tr>
<td>tz</td>
<td>lower case time-zone abbreviation (only supported in to_char)</td>
</tr>
<tr>
<td>TZH</td>
<td>time-zone hours</td>
</tr>
<tr>
<td>TZM</td>
<td>time-zone minutes</td>
</tr>
<tr>
<td>OF</td>
<td>time-zone offset from UTC (only supported in to_char)</td>
</tr>
</tbody>
</table>

Modifiers can be applied to any template pattern to alter its behavior. For example, FMMonth is the Month pattern with the FM modifier. Table 9.25 shows the modifier patterns for date/time formatting.

**Table 9.25. Template Pattern Modifiers for Date/Time Formatting**

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>FM prefix</td>
<td>fill mode (suppress leading zeroes and padding blanks)</td>
<td>FMMonth</td>
</tr>
<tr>
<td>TH suffix</td>
<td>upper case ordinal number suffix</td>
<td>DDTH, e.g., 12TH</td>
</tr>
<tr>
<td>th suffix</td>
<td>lower case ordinal number suffix</td>
<td>DDth, e.g., 12th</td>
</tr>
<tr>
<td>FX prefix</td>
<td>fixed format global option (see usage notes)</td>
<td>FX Month DD Day</td>
</tr>
<tr>
<td>TM prefix</td>
<td>translation mode (print localized day and month names based on lc_time)</td>
<td>TMMonth</td>
</tr>
<tr>
<td>SP suffix</td>
<td>spell mode (not implemented)</td>
<td>DDSP</td>
</tr>
</tbody>
</table>

Usage notes for date/time formatting:

- **FM** suppresses leading zeroes and trailing blanks that would otherwise be added to make the output of a pattern be fixed-width. In PostgreSQL, FM modifies only the next specification, while in Oracle FM affects all subsequent specifications, and repeated FM modifiers toggle fill mode on and off.

- **TM** does not include trailing blanks. `to_timestamp` and `to_date` ignore the TM modifier.

- `to_timestamp` and `to_date` skip multiple blank spaces in the input string unless the FX option is used. For example, `to_timestamp('2000 JUN', 'YYYY MON')` works, but `to_timestamp('2000 JUN', 'FXYYYY MON')` returns an error because `to_timestamp` expects one space only. FX must be specified as the first item in the template.

- Ordinary text is allowed in `to_char` templates and will be output literally. You can put a substring in double quotes to force it to be interpreted as literal text even if it contains template patterns. For example, in ""Hello Year "YYYY"", the YYYY will be replaced by the year data, but the single Y in Year will not be. In `to_date`, `to_number`, and `to_timestamp`, literal text and double-quoted strings result in skipping the number of characters contained in the string; for example "XX" skips two input characters (whether or not they are XX).

- If you want to have a double quote in the output you must precede it with a backslash, for example '""YYYY Month""'. Backslashes are not otherwise special outside of double-quoted strings.
Within a double-quoted string, a backslash causes the next character to be taken literally, whatever it is (but this has no special effect unless the next character is a double quote or another backslash).

**In `to_timestamp` and `to_date`,** if the year format specification is less than four digits, e.g. `YYY`, and the supplied year is less than four digits, the year will be adjusted to be nearest to the year 2020, e.g. 95 becomes 1995.

**In `to_timestamp` and `to_date`,** the `YYYY` conversion has a restriction when processing years with more than 4 digits. You must use some non-digit character or template after `YYYY`, otherwise the year is always interpreted as 4 digits. For example (with the year 20000): `to_date('200001131', 'YYYYMMDD')` will be interpreted as a 4-digit year; instead use a non-digit separator after the year, like `to_date('20000-1131', 'YYYY-MMDD')` or `to_date('20000Nov31', 'YYYYMonDD')`.

**In `to_timestamp` and `to_date`,** the `CC` (century) field is accepted but ignored if there is a `YYY, YYYY` or `Y, YYYY` field. If `CC` is used with `YY` or `Y` then the result is computed as that year in the specified century. If the century is specified but the year is not, the first year of the century is assumed.

**In `to_timestamp` and `to_date`,** weekday names or numbers (`DAY`, `D`, and related field types) are accepted but are ignored for purposes of computing the result. The same is true for quarter (`Q`) fields.

**In `to_timestamp` and `to_date`,** an ISO 8601 week-numbering date (as distinct from a Gregorian date) can be specified in one of two ways:

- Year, week number, and weekday: for example `to_date('2006-42-4', 'IYYY-IW-ID')` returns the date 2006-10-19. If you omit the weekday it is assumed to be 1 (Monday).
- Year and day of year: for example `to_date('2006-291', 'IYYY-IDDD')` also returns 2006-10-19.

Attempting to enter a date using a mixture of ISO 8601 week-numbering fields and Gregorian date fields is nonsensical, and will cause an error. In the context of an ISO 8601 week-numbering year, the concept of a “month” or “day of month” has no meaning. In the context of a Gregorian year, the ISO week has no meaning.

**Caution**

While `to_date` will reject a mixture of Gregorian and ISO week-numbering date fields, `to_char` will not, since output format specifications like `YYYY-MM-DD (IYYY-IDDD)` can be useful. But avoid writing something like `IYYY-MM-DD;` that would yield surprising results near the start of the year. (See Section 9.9.1 for more information.)

**In `to_timestamp`,** millisecond (`MS`) or microsecond (`US`) fields are used as the seconds digits after the decimal point. For example `to_timestamp('12.3', 'SS.MS')` is not 3 milliseconds, but 300, because the conversion treats it as 12 + 0.3 seconds. So, for the format `SS.MS`, the input values 12.3, 12.30, and 12.300 specify the same number of milliseconds. To get three milliseconds, one must write 12.003, which the conversion treats as 12 + 0.003 = 12.003 seconds.

Here is a more complex example: `to_timestamp('15:12:02.020.001230', 'HH24:MI:SS.MS.US')` is 15 hours, 12 minutes, and 2 seconds + 20 milliseconds + 1230 microseconds = 2.021230 seconds.

**`to_char(..., 'ID')`’s day of the week numbering matches the `extract(isodow from ...)` function, but `to_char(..., 'D')`’s does not match `extract(dow from ...)`’s day numbering.**
Functions and Operators

- `to_char(interval)` formats `HH` and `HH12` as shown on a 12-hour clock, for example zero hours and 36 hours both output as 12, while `HH24` outputs the full hour value, which can exceed 23 in an interval value.

Table 9.26 shows the template patterns available for formatting numeric values.

**Table 9.26. Template Patterns for Numeric Formatting**

<table>
<thead>
<tr>
<th>Pattern</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>9</td>
<td>digit position (can be dropped if insignificant)</td>
</tr>
<tr>
<td>0</td>
<td>digit position (will not be dropped, even if insignificant)</td>
</tr>
<tr>
<td>. (period)</td>
<td>decimal point</td>
</tr>
<tr>
<td>, (comma)</td>
<td>group (thousands) separator</td>
</tr>
<tr>
<td>PR</td>
<td>negative value in angle brackets</td>
</tr>
<tr>
<td>S</td>
<td>sign anchored to number (uses locale)</td>
</tr>
<tr>
<td>L</td>
<td>currency symbol (uses locale)</td>
</tr>
<tr>
<td>D</td>
<td>decimal point (uses locale)</td>
</tr>
<tr>
<td>G</td>
<td>group separator (uses locale)</td>
</tr>
<tr>
<td>MI</td>
<td>minus sign in specified position (if number &lt; 0)</td>
</tr>
<tr>
<td>PL</td>
<td>plus sign in specified position (if number &gt; 0)</td>
</tr>
<tr>
<td>SG</td>
<td>plus/minus sign in specified position</td>
</tr>
<tr>
<td>RN</td>
<td>Roman numeral (input between 1 and 3999)</td>
</tr>
<tr>
<td>TH or th</td>
<td>ordinal number suffix</td>
</tr>
<tr>
<td>V</td>
<td>shift specified number of digits (see notes)</td>
</tr>
<tr>
<td>EEEEE</td>
<td>exponent for scientific notation</td>
</tr>
</tbody>
</table>

Usage notes for numeric formatting:

- 0 specifies a digit position that will always be printed, even if it contains a leading/trailing zero. 9 also specifies a digit position, but if it is a leading zero then it will be replaced by a space, while if it is a trailing zero and fill mode is specified then it will be deleted. (For `to_number()`, these two pattern characters are equivalent.)

- The pattern characters `S`, `L`, `D`, and `G` represent the sign, currency symbol, decimal point, and thousands separator characters defined by the current locale (see `lc_monetary` and `lc_numeric`). The pattern characters period and comma represent those exact characters, with the meanings of decimal point and thousands separator, regardless of locale.

- If no explicit provision is made for a sign in `to_char()`’s pattern, one column will be reserved for the sign, and it will be anchored to (appear just left of) the number. If S appears just left of some 9’s, it will likewise be anchored to the number.

- A sign formatted using `SG`, `PL`, or `MI` is not anchored to the number; for example, `to_char(-12, 'MI9999')` produces `- 12` but `to_char(-12, '99999')` produces `-12`. (The Oracle implementation does not allow the use of `MI` before 9, but rather requires that 9 precede `MI`.)

- `TH` does not convert values less than zero and does not convert fractional numbers.

- `PL`, `SG`, and `TH` are PostgreSQL extensions.

- In `to_number`, if non-data template patterns such as `L` or `TH` are used, the corresponding number of input characters are skipped, whether or not they match the template pattern, unless they are data characters (that is, digits, sign, decimal point, or comma). For example, `TH` would skip two non-data characters.
• $V$ with `to_char` multiplies the input values by $10^n$, where $n$ is the number of digits following $V$. $V$ with `to_number` divides in a similar manner. `to_char` and `to_number` do not support the use of $V$ combined with a decimal point (e.g., `99.9V99` is not allowed).

• `EEEE` (scientific notation) cannot be used in combination with any of the other formatting patterns or modifiers other than digit and decimal point patterns, and must be at the end of the format string (e.g., `9.99EEEE` is a valid pattern).

Certain modifiers can be applied to any template pattern to alter its behavior. For example, `FM99.99` is the `99.99` pattern with the `FM` modifier. Table 9.27 shows the modifier patterns for numeric formatting.

**Table 9.27. Template Pattern Modifiers for Numeric Formatting**

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>FM prefix</td>
<td>fill mode (suppress trailing zeroes and padding blanks)</td>
<td><code>FM99.99</code></td>
</tr>
<tr>
<td>TH suffix</td>
<td>upper case ordinal number suffix</td>
<td><code>999TH</code></td>
</tr>
<tr>
<td>th suffix</td>
<td>lower case ordinal number suffix</td>
<td><code>999th</code></td>
</tr>
</tbody>
</table>

Table 9.28 shows some examples of the use of the `to_char` function.

**Table 9.28. `to_char` Examples**

<table>
<thead>
<tr>
<th>Expression</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_char(current_timestamp, 'Day, DD HH12:MI:SS')</code></td>
<td>'Tuesday, 06 05:39:18'</td>
</tr>
<tr>
<td><code>to_char(current_timestamp, 'FM-Day, FMDD HH12:MI:SS')</code></td>
<td>'Tuesday, 6 05:39:18'</td>
</tr>
<tr>
<td><code>to_char(-0.1, '99.99')</code></td>
<td>'-.10'</td>
</tr>
<tr>
<td><code>to_char(-0.1, 'FM9.99')</code></td>
<td>'-.1'</td>
</tr>
<tr>
<td><code>to_char(-0.1, 'FM90.99')</code></td>
<td>'-0.1'</td>
</tr>
<tr>
<td><code>to_char(0.1, '0.9')</code></td>
<td>'.01'</td>
</tr>
<tr>
<td><code>to_char(12, '9990999.9')</code></td>
<td>'0012.0'</td>
</tr>
<tr>
<td><code>to_char(12, 'FM9990999.9')</code></td>
<td>'0012.'</td>
</tr>
<tr>
<td><code>to_char(485, '999')</code></td>
<td>'485'</td>
</tr>
<tr>
<td><code>to_char(-485, '999')</code></td>
<td>'-485'</td>
</tr>
<tr>
<td><code>to_char(485, '9 9 9')</code></td>
<td>'485'</td>
</tr>
<tr>
<td><code>to_char(1485, '9 999')</code></td>
<td>'1485'</td>
</tr>
<tr>
<td><code>to_char(1485, '9G999')</code></td>
<td>'1 485'</td>
</tr>
<tr>
<td><code>to_char(148.5, '999.999')</code></td>
<td>'148.500'</td>
</tr>
<tr>
<td><code>to_char(148.5, 'FM999.999')</code></td>
<td>'148.5'</td>
</tr>
<tr>
<td><code>to_char(148.5, 'FM999.990')</code></td>
<td>'148.500'</td>
</tr>
<tr>
<td><code>to_char(148.5, '999D999')</code></td>
<td>'148,500'</td>
</tr>
<tr>
<td><code>to_char(3148.5, '9G999D999')</code></td>
<td>'3 148,500'</td>
</tr>
<tr>
<td><code>to_char(-485, '999S')</code></td>
<td>'-485'</td>
</tr>
<tr>
<td><code>to_char(-485, '999MI')</code></td>
<td>'-485-'</td>
</tr>
<tr>
<td><code>to_char(485, '999MI')</code></td>
<td>'485'</td>
</tr>
<tr>
<td><code>to_char(485, 'FM999MI')</code></td>
<td>'485'</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Expression</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_char(485, 'PL999')</code></td>
<td>'+485'</td>
</tr>
<tr>
<td><code>to_char(485, 'SG999')</code></td>
<td>'+485'</td>
</tr>
<tr>
<td><code>to_char(-485, 'SG999')</code></td>
<td>'-485'</td>
</tr>
<tr>
<td><code>to_char(-485, '9SG99')</code></td>
<td>'-485&gt;'</td>
</tr>
<tr>
<td><code>to_char(485, '99PR')</code></td>
<td>'485'</td>
</tr>
<tr>
<td><code>to_char(485, 'L999')</code></td>
<td>'DM 485'</td>
</tr>
<tr>
<td><code>to_char(485, 'RN')</code></td>
<td>'CDLXXXV'</td>
</tr>
<tr>
<td><code>to_char(485, 'FMRN')</code></td>
<td>'CDLXXXV'</td>
</tr>
<tr>
<td><code>to_char(5.2, 'FMRN')</code></td>
<td>'V'</td>
</tr>
<tr>
<td><code>to_char(482, '999th')</code></td>
<td>'482nd'</td>
</tr>
<tr>
<td><code>to_char(485, &quot;Good number:&quot;999')</code></td>
<td>'Good number: 485'</td>
</tr>
<tr>
<td><code>to_char(485.8, &quot;Pre:&quot;999 Post:&quot; .999')</code></td>
<td>'Pre: 485 Post: .800'</td>
</tr>
<tr>
<td><code>to_char(12, '99V999')</code></td>
<td>'12000'</td>
</tr>
<tr>
<td><code>to_char(12.4, '99V999')</code></td>
<td>'12400'</td>
</tr>
<tr>
<td><code>to_char(12.45, '99V9')</code></td>
<td>'125'</td>
</tr>
<tr>
<td><code>to_char(0.0004859, '9.99EEEE')</code></td>
<td>'4.86e-04'</td>
</tr>
</tbody>
</table>

### 9.9. Date/Time Functions and Operators

Table 9.30 shows the available functions for date/time value processing, with details appearing in the following subsections. Table 9.29 illustrates the behaviors of the basic arithmetic operators (+, *, etc.). For formatting functions, refer to Section 9.8. You should be familiar with the background information on date/time data types from Section 8.5.

All the functions and operators described below that take time or timestamp inputs actually come in two variants: one that takes time with time zone or timestamp with time zone, and one that takes time without time zone or timestamp without time zone. For brevity, these variants are not shown separately. Also, the + and * operators come in commutative pairs (for example both date + integer and integer + date); we show only one of each such pair.

#### Table 9.29. Date/Time Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>date '2001-09-28' + integer '7'</td>
<td>date '2001-10-05'</td>
</tr>
<tr>
<td>+</td>
<td>date '2001-09-28' + interval '1 hour'</td>
<td>timestamp '2001-09-28 01:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>date '2001-09-28' + time '03:00'</td>
<td>timestamp '2001-09-28 03:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>interval '1 day' + interval '1 hour'</td>
<td>interval '1 day 01:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>timestamp '2001-09-28 01:00' + interval '23 hours'</td>
<td>timestamp '2001-09-28 00:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>time '01:00' + interval '3 hours'</td>
<td>time '04:00:00'</td>
</tr>
<tr>
<td>Operator</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>---------</td>
<td>--------</td>
</tr>
<tr>
<td>-</td>
<td>- interval '23 hours'</td>
<td>interval '-23:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-10-01' - date '2001-09-28'</td>
<td>integer '3' (days)</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-10-01' - integer '7'</td>
<td>date '2001-09-24'</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-09-28' - interval '1 hour'</td>
<td>timestamp '2001-09-27 23:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>time '05:00' - time '03:00'</td>
<td>interval '02:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>time '05:00' - interval '2 hours'</td>
<td>time '03:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>timestamp '2001-09-28 23:00' - interval '23 hours'</td>
<td>timestamp '2001-09-28 00:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>interval '1 day' - interval '1 hour'</td>
<td>interval '1 day 01:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>timestamp '2001-09-29 03:00' - timestamp '2001-09-27 12:00'</td>
<td>interval '1 day 15:00:00'</td>
</tr>
<tr>
<td>*</td>
<td>900 * interval '1 second'</td>
<td>interval '00:15:00'</td>
</tr>
<tr>
<td>*</td>
<td>21 * interval '1 day'</td>
<td>interval '21 days'</td>
</tr>
<tr>
<td>*</td>
<td>double precision '3.5' * interval '1 hour'</td>
<td>interval '03:30:00'</td>
</tr>
<tr>
<td>/</td>
<td>interval '1 hour' / double precision '1.5'</td>
<td>interval '00:40:00'</td>
</tr>
</tbody>
</table>

Table 9.30. Date/Time Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>age(time-stamp, time-stamp)</td>
<td>interval</td>
<td>Subtract arguments, producing a &quot;symbolic&quot; result that uses years and months, rather than just days</td>
<td>age(time-stamp '2001-04-10', time-stamp '1957-06-13')</td>
<td>43 years 9 mons 27 days</td>
</tr>
<tr>
<td>age(time-stamp)</td>
<td>interval</td>
<td>Subtract from current_date (at midnight)</td>
<td>age(time-stamp '1957-06-13')</td>
<td>43 years 8 mons 3 days</td>
</tr>
<tr>
<td>clock_time-stamp() with time zone</td>
<td>timestamp</td>
<td>Current date and time (changes during statement execution); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>current_date</td>
<td>date</td>
<td>Current date; see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>current_time time with time zone</td>
<td>Current time of day; see Section 9.9.4</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>-------------------------</td>
<td>-------------</td>
<td>-------------------------------------------------------</td>
<td>----------------------------------------------</td>
<td>---------------</td>
</tr>
<tr>
<td>current_timestamp</td>
<td>timestamp</td>
<td>Current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>t_timestamp</td>
<td>time zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>date_part(text,timestamp)</td>
<td>double precision</td>
<td>Get subfield (equivalent to extract); see Section 9.9.1</td>
<td>date_part('hour', timestamp '2001-02-16 20:38:40')</td>
<td>20</td>
</tr>
<tr>
<td>date_part(text, interval)</td>
<td>double precision</td>
<td>Get subfield (equivalent to extract); see Section 9.9.1</td>
<td>date_part('month', interval '2 years 3 months')</td>
<td>3</td>
</tr>
<tr>
<td>date_trunc(</td>
<td>timestamp</td>
<td>Truncate to specified precision; see also Section 9.9.2</td>
<td>date_trunc('day', timestamp '2001-02-16 20:00:00')</td>
<td>2001-02-16 20:00:00</td>
</tr>
<tr>
<td>interval)</td>
<td></td>
<td></td>
<td>date_trunc('hour', interval '2 03:00:00 days 3 hours 40 minutes')</td>
<td></td>
</tr>
<tr>
<td>extract(</td>
<td>double precision</td>
<td>Get subfield; see Section 9.9.1</td>
<td>extract(hour from timestamp '2001-02-16 20:38:40')</td>
<td>20</td>
</tr>
<tr>
<td>field, from sion timestamp)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>extract(</td>
<td>double precision</td>
<td>Get subfield; see Section 9.9.1</td>
<td>extract(month from interval '2 years 3 months')</td>
<td>3</td>
</tr>
<tr>
<td>field, from sion interval)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>isfinite(date)</td>
<td>boolean</td>
<td>Test for finite date (not +/-infinity)</td>
<td>isfinite(date '2001-02-16')</td>
<td>true</td>
</tr>
<tr>
<td>isfinite(time- stamp)</td>
<td>boolean</td>
<td>Test for finite time stamp (not +/-infinity)</td>
<td>isfinite(time- stamp '2001-02-16 21:28:30')</td>
<td>true</td>
</tr>
<tr>
<td>isfinite(interval)</td>
<td>boolean</td>
<td>Test for finite interval</td>
<td>isfinite(interval '4 hours')</td>
<td></td>
</tr>
<tr>
<td>justify_days(interval)</td>
<td>interval</td>
<td>Adjust interval so 30-day time periods are represented as months</td>
<td>justify_days(interval '35 days')</td>
<td>1 mon 5 days</td>
</tr>
<tr>
<td>justify_hours(interval)</td>
<td>interval</td>
<td>Adjust interval so 24-hour time periods are represented as days</td>
<td>justify_hours(interval '27 hours')</td>
<td>1 day</td>
</tr>
<tr>
<td>justify_in_interval(interval)</td>
<td>interval</td>
<td>Adjust interval using justify_in_interval('29 days') and fy_days</td>
<td>justify_in_interval(interval '23:00:00')</td>
<td>29</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>----------------------------------</td>
<td>-------------------------------</td>
</tr>
<tr>
<td>justify_hours</td>
<td></td>
<td>adjustment of a given interval by an additional sign adjustment</td>
<td>'1 month -1 hour'</td>
<td></td>
</tr>
<tr>
<td>localtime</td>
<td>time</td>
<td>current time of day; see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>localtime-stamp</td>
<td>timestamp</td>
<td>current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>make_date(year, month, day)</td>
<td>date</td>
<td>create date from year, month and day fields</td>
<td>2013, 7, 15</td>
<td>2013-07-15</td>
</tr>
<tr>
<td>make_interval(year, month, week, day, hour, min, sec)</td>
<td>interval</td>
<td>create interval from years, months, weeks, days, hours, minutes and seconds fields</td>
<td>days =&gt; 10</td>
<td>10 days</td>
</tr>
<tr>
<td>make_time(hour, min, sec)</td>
<td>time</td>
<td>create time from hour, minute and seconds fields</td>
<td>8, 15, 23.5</td>
<td>08:15:23.5</td>
</tr>
<tr>
<td>make_time-stamp(year, month, day, hour, min, sec)</td>
<td>timestamp</td>
<td>create timestamp from year, month, day, hour, minute and seconds fields</td>
<td>2013, 7, 15, 8, 15, 23.5</td>
<td>2013-07-15 08:15:23.5</td>
</tr>
<tr>
<td>make_time-stamptz(year, month, day, hour, min, sec, time zone)</td>
<td>timestamp with time zone</td>
<td>create timestamp with time zone from year, month, day, hour, minute, and seconds fields; if timezone is not specified, the current time zone is used</td>
<td>2013, 7, 15, 8, 15, 23.5+01</td>
<td>2013-07-15 08:15:23.5+01</td>
</tr>
</tbody>
</table>

247
## Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>now()</td>
<td>timestamp with time zone</td>
<td>Current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>statement_timestamp()</td>
<td>timestamp with time zone</td>
<td>Current date and time (start of current statement); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>to_timestamp(double precision)</td>
<td>timestamp with time zone</td>
<td>Convert Unix epoch (seconds since 1970-01-01 00:00:00+00) to timestamp</td>
<td>to_timestamp(1284352323)</td>
<td>2010-09-13 04:32:03+00</td>
</tr>
</tbody>
</table>

In addition to these functions, the SQL `OVERLAPS` operator is supported:

```
(start1, end1) OVERLAPS (start2, end2)
(start1, length1) OVERLAPS (start2, length2)
```

This expression yields true when two time periods (defined by their endpoints) overlap, false when they do not overlap. The endpoints can be specified as pairs of dates, times, or time stamps; or as a date, time, or time stamp followed by an interval. When a pair of values is provided, either the start or the end can be written first; `OVERLAPS` automatically takes the earlier value of the pair as the start. Each time period is considered to represent the half-open interval `start <= time < end`, unless `start` and `end` are equal in which case it represents that single time instant. This means for instance that two time periods with only an endpoint in common do not overlap.

```
SELECT (DATE '2001-02-16', DATE '2001-12-21') OVERLAPS (DATE '2001-10-30', DATE '2002-10-30');
Result: true
```
```
SELECT (DATE '2001-02-16', INTERVAL '100 days') OVERLAPS (DATE '2001-10-30', DATE '2002-10-30');
Result: false
```
```
SELECT (DATE '2001-10-29', DATE '2001-10-30') OVERLAPS (DATE '2001-10-30', DATE '2001-10-31');
Result: false
```
```
SELECT (DATE '2001-10-30', DATE '2001-10-30') OVERLAPS (DATE '2001-10-30', DATE '2001-10-31');
Result: true
```

When adding an interval value to (or subtracting an interval value from) a timestamp with time zone value, the days component advances or decrements the date of the timestamp with time zone by the indicated number of days. Across daylight saving time changes (when the session time zone is set to a time zone that recognizes DST), this means interval '1 day'
does not necessarily equal interval '24 hours'. For example, with the session time zone set to CST7CDT, timestamp with time zone '2005-04-02 12:00-07' + interval '1 day' will produce timestamp with time zone '2005-04-03 12:00-06', while adding interval '24 hours' to the same initial timestamp with time zone produces timestamp with time zone '2005-04-03 13:00-06', as there is a change in daylight saving time at 2005-04-03 02:00 in time zone CST7CDT.

Note there can be ambiguity in the months field returned by age because different months have different numbers of days. PostgreSQL's approach uses the month from the earlier of the two dates when calculating partial months. For example, age('2004-06-01', '2004-04-30') uses April to yield 1 mon 1 day, while using May would yield 1 mon 2 days because May has 31 days, while April has only 30.

Subtraction of dates and timestamps can also be complex. One conceptually simple way to perform subtraction is to convert each value to a number of seconds using EXTRACT(EPOCH FROM ...), then subtract the results; this produces the number of seconds between the two values. This will adjust for the number of days in each month, timezone changes, and daylight saving time adjustments. Subtraction of date or timestamp values with the "-" operator returns the number of days (24-hours) and hours/minutes/seconds between the values, making the same adjustments. The age function returns years, months, days, and hours/minutes/seconds, performing field-by-field subtraction and then adjusting for negative field values. The following queries illustrate the differences in these approaches. The sample results were produced with timezone = 'US/Eastern'; there is a daylight saving time change between the two dates used:

```sql
SELECT EXTRACT(EPOCH FROM timestamptz '2013-07-01 12:00:00') - EXTRACT(EPOCH FROM timestamptz '2013-03-01 12:00:00');
Result: 10537200
SELECT (EXTRACT(EPOCH FROM timestamptz '2013-07-01 12:00:00') - EXTRACT(EPOCH FROM timestamptz '2013-03-01 12:00:00')) / 60 / 60 / 24;
Result: 121.958333333333
SELECT timestamptz '2013-07-01 12:00:00' - timestamptz '2013-03-01 12:00:00';
Result: 121 days 23:00:00
SELECT age(timestamptz '2013-07-01 12:00:00', timestamptz '2013-03-01 12:00:00');
Result: 4 mons
```

### 9.9.1. EXTRACT, date_part

**EXTRACT(field FROM source)**

The extract function retrieves subfields such as year or hour from date/time values. source must be a value expression of type timestamp, time, or interval. (Expressions of type date are cast to timestamp and can therefore be used as well.) field is an identifier or string that selects what field to extract from the source value. The extract function returns values of type double precision. The following are valid field names:

- **century**
  
  The century

```sql
SELECT EXTRACT(CENTURY FROM TIMESTAMP '2000-12-16 12:21:13');
Result: 20
SELECT EXTRACT(CENTURY FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 21
```
The first century starts at 0001-01-01 00:00:00 AD, although they did not know it at the time. This definition applies to all Gregorian calendar countries. There is no century number 0, you go from -1 century to 1 century. If you disagree with this, please write your complaint to: Pope, Cathedral Saint-Peter of Roma, Vatican.

day

For timestamp values, the day (of the month) field (1 - 31); for interval values, the number of days

SELECT EXTRACT(DAY FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 16

SELECT EXTRACT(DAY FROM INTERVAL '40 days 1 minute');
Result: 40

decade

The year field divided by 10

SELECT EXTRACT(DECADE FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 200

dow

The day of the week as Sunday (0) to Saturday (6)

SELECT EXTRACT(DOW FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 5

Note that extract's day of the week numbering differs from that of the to_char(..., 'D') function.

doy

The day of the year (1 - 365/366)

SELECT EXTRACT(DOY FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 47

epoch

For timestamp with time zone values, the number of seconds since 1970-01-01 00:00:00 UTC (can be negative); for date and timestamp values, the number of seconds since 1970-01-01 00:00:00 local time; for interval values, the total number of seconds in the interval

SELECT EXTRACT(EPOCH FROM TIMESTAMP WITH TIME ZONE '2001-02-16 20:38:40.12-08');
Result: 982384720.12

SELECT EXTRACT(EPOCH FROM INTERVAL '5 days 3 hours');
Result: 442800

You can convert an epoch value back to a time stamp with to_timestamp:

SELECT to_timestamp(982384720.12);
Functions and Operators

Result: 2001-02-17 04:38:40.12+00

hour

The hour field (0 - 23)

SELECT EXTRACT(HOUR FROM TIMESTAMPTZ '2001-02-16 20:38:40');
Result: 20

isodow

The day of the week as Monday (1) to Sunday (7)

SELECT EXTRACT(ISODOW FROM TIMESTAMPTZ '2001-02-18 20:38:40');
Result: 7

This is identical to dow except for Sunday. This matches the ISO 8601 day of the week numbering.

isoyear

The ISO 8601 week-numbering year that the date falls in (not applicable to intervals)

SELECT EXTRACT(ISOYEAR FROM DATE '2006-01-01');
Result: 2005
SELECT EXTRACT(ISOYEAR FROM DATE '2006-01-02');
Result: 2006

Each ISO 8601 week-numbering year begins with the Monday of the week containing the 4th of January, so in early January or late December the ISO year may be different from the Gregorian year. See the week field for more information.

This field is not available in PostgreSQL releases prior to 8.3.

microseconds

The seconds field, including fractional parts, multiplied by 1 000 000; note that this includes full seconds

SELECT EXTRACT(MICROSECONDS FROM TIMESTAMPTZ '2001-02-16 20:38:40');
Result: 28500000

millennium

The millennium

SELECT EXTRACT(MILLENNIUM FROM TIMESTAMPTZ '2001-02-16 20:38:40');
Result: 3

Years in the 1900s are in the second millennium. The third millennium started January 1, 2001.

milliseconds

The seconds field, including fractional parts, multiplied by 1000. Note that this includes full seconds.

SELECT EXTRACT(MILLISECONDS FROM TIMESTAMPTZ '2001-02-16 20:38:40');
Result: 28500
Functions and Operators

minute

The minutes field (0 - 59)

```
SELECT EXTRACT(MINUTE FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 38
```

month

For timestamp values, the number of the month within the year (1 - 12); for interval values, the number of months, modulo 12 (0 - 11)

```
SELECT EXTRACT(MONTH FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 2

SELECT EXTRACT(MONTH FROM INTERVAL '2 years 3 months');
Result: 3

SELECT EXTRACT(MONTH FROM INTERVAL '2 years 13 months');
Result: 1
```

quarter

The quarter of the year (1 - 4) that the date is in

```
SELECT EXTRACT(QUARTER FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 1
```

second

The seconds field, including fractional parts (0 - 59\(^1\))

```
SELECT EXTRACT(SECOND FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 40

SELECT EXTRACT(SECOND FROM TIME '17:12:28.5');
Result: 28.5
```

timezone

The time zone offset from UTC, measured in seconds. Positive values correspond to time zones east of UTC, negative values to zones west of UTC. (Technically, PostgreSQL does not use UTC because leap seconds are not handled.)

timezone_hour

The hour component of the time zone offset

timezone_minute

The minute component of the time zone offset

week

The number of the ISO 8601 week-numbering week of the year. By definition, ISO weeks start on Mondays and the first week of a year contains January 4 of that year. In other words, the first Thursday of a year is in week 1 of that year.

\(^1\)60 if leap seconds are implemented by the operating system
In the ISO week-numbering system, it is possible for early-January dates to be part of the 52nd or 53rd week of the previous year, and for late-December dates to be part of the first week of the next year. For example, 2005-01-01 is part of the 53rd week of year 2004, and 2006-01-01 is part of the 52nd week of year 2005, while 2012-12-31 is part of the first week of 2013. It's recommended to use the isoyear field together with week to get consistent results.

```sql
SELECT EXTRACT(WEEK FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 7
```

**year**

The year field. Keep in mind there is no 0 AD, so subtracting BC years from AD years should be done with care.

```sql
SELECT EXTRACT(YEAR FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 2001
```

**Note**

When the input value is +/-Infinity, extract returns +/-Infinity for monotonically-increasing fields (epoch, julian, year, isoyear, decade, century, and millennium). For other fields, NULL is returned. PostgreSQL versions before 9.6 returned zero for all cases of infinite input.

The `extract` function is primarily intended for computational processing. For formatting date/time values for display, see Section 9.8.

The `date_part` function is modeled on the traditional Ingres equivalent to the SQL-standard function `extract`:

```sql
date_part('field', source)
```

Note that here the `field` parameter needs to be a string value, not a name. The valid field names for `date_part` are the same as for `extract`.

```sql
SELECT date_part('day', TIMESTAMP '2001-02-16 20:38:40');
Result: 16
```

```sql
SELECT date_part('hour', INTERVAL '4 hours 3 minutes');
Result: 4
```

### 9.9.2. `date_trunc`

The function `date_trunc` is conceptually similar to the `trunc` function for numbers.

```sql
date_trunc('field', source)
```

`source` is a value expression of type `timestamp` or `interval`. (Values of type `date` and `time` are cast automatically to `timestamp` or `interval`, respectively.) `field` selects to which precision to truncate the input value. The return value is of type `timestamp` or `interval` with all fields that are less significant than the selected one set to zero (or one, for day and month).
Valid values for field are:

- microseconds
- milliseconds
- second
- minute
- hour
- day
- week
- month
- quarter
- year
- decade
- century
- millennium

Examples:

```sql
SELECT date_trunc('hour', TIMESTAMP '2001-02-16 20:38:40');
Result: 2001-02-16 20:00:00

SELECT date_trunc('year', TIMESTAMP '2001-02-16 20:38:40');
Result: 2001-01-01 00:00:00
```

### 9.9.3. AT TIME ZONE

The AT TIME ZONE converts time stamp without time zone to/from time stamp with time zone, and time values to different time zones. Table 9.31 shows its variants.

#### Table 9.31. AT TIME ZONE Variants

<table>
<thead>
<tr>
<th>Expression</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>timestamp without time zone AT TIME ZONE zone</td>
<td>timestamp with time zone</td>
<td>Treat given time stamp without time zone as located in the specified time zone</td>
</tr>
<tr>
<td>timestamp with time zone AT TIME ZONE zone</td>
<td>timestamp without time zone</td>
<td>Convert given time stamp with time zone to the new time zone, with no time zone designation</td>
</tr>
<tr>
<td>time with time zone AT TIME ZONE zone</td>
<td>time with time zone</td>
<td>Convert given time with time zone to the new time zone</td>
</tr>
</tbody>
</table>

In these expressions, the desired time zone name can be specified either as a text string (e.g., 'America/Los_Angeles') or as an interval (e.g., INTERVAL '-08:00'). In the text case, a time zone name can be specified in any of the ways described in Section 8.5.3.

Examples (assuming the local time zone is America/Los_Angeles):

```sql
SELECT TIMESTAMP '2001-02-16 20:38:40' AT TIME ZONE 'America/Denver';
Result: 2001-02-16 19:38:40-08

SELECT TIMESTAMP WITH TIME ZONE '2001-02-16 20:38:40-05' AT TIME ZONE 'America/Denver';
Result: 2001-02-16 18:38:40
```
The first example adds a time zone to a value that lacks it, and displays the value using the current TimeZone setting. The second example shifts the time stamp with time zone value to the specified time zone, and returns the value without a time zone. This allows storage and display of values different from the current TimeZone setting. The third example converts Tokyo time to Chicago time. Converting time values to other time zones uses the currently active time zone rules since no date is supplied.

The function timezone (zone, timestamp) is equivalent to the SQL-conforming construct timestamp AT TIME ZONE zone.

### 9.9.4. Current Date/Time

PostgreSQL provides a number of functions that return values related to the current date and time. These SQL-standard functions all return values based on the start time of the current transaction:

- `CURRENT_DATE`
- `CURRENT_TIME`
- `CURRENT_TIMESTAMP`
- `CURRENT_TIME(precision)`
- `CURRENT_TIMESTAMP(precision)`
- `LOCALTIME`
- `LOCALTIMESTAMP`
- `LOCALTIME(precision)`
- `LOCALTIMESTAMP(precision)`

`CURRENT_TIME` and `CURRENT_TIMESTAMP` deliver values with time zone; `LOCALTIME` and `LOCALTIMESTAMP` deliver values without time zone.

`CURRENT_TIME, CURRENT_TIMESTAMP, LOCALTIME, LOCALTIMESTAMP` can optionally take a precision parameter, which causes the result to be rounded to that many fractional digits in the seconds field. Without a precision parameter, the result is given to the full available precision.

Some examples:

```
SELECT CURRENT_TIME;
Result: 14:39:53.662522-05
```

```
SELECT CURRENT_DATE;
Result: 2001-12-23
```

```
SELECT CURRENT_TIMESTAMP;
Result: 2001-12-23 14:39:53.662522-05
```

```
SELECT CURRENT_TIMESTAMP(2);
Result: 2001-12-23 14:39:53.66-05
```

```
SELECT LOCALTIMESTAMP;
Result: 2001-12-23 14:39:53.662522
```

Since these functions return the start time of the current transaction, their values do not change during the transaction. This is considered a feature: the intent is to allow a single transaction to have a consistent notion of the “current” time, so that multiple modifications within the same transaction bear the same time stamp.
PostgreSQL also provides functions that return the start time of the current statement, as well as the actual current time at the instant the function is called. The complete list of non-SQL-standard time functions is:

transaction_timestamp()
statement_timestamp()
clock_timestamp()
timeofday()
now()

transaction_timestamp() is equivalent to CURRENT_TIMESTAMP, but is named to clearly reflect what it returns. statement_timestamp() returns the start time of the current statement (more specifically, the time of receipt of the latest command message from the client). statement_timestamp() and transaction_timestamp() return the same value during the first command of a transaction, but might differ during subsequent commands. clock_timestamp() returns the actual current time, and therefore its value changes even within a single SQL command. timeofday() is a historical PostgreSQL function. Like clock_timestamp(), it returns the actual current time, but as a formatted text string rather than a timestamp with time zone value. now() is a traditional PostgreSQL equivalent to transaction_timestamp().

All the date/time data types also accept the special literal value now to specify the current date and time (again, interpreted as the transaction start time). Thus, the following three all return the same result:

```
SELECT CURRENT_TIMESTAMP;
SELECT now();
SELECT TIMESTAMP 'now';  -- incorrect for use with DEFAULT
```

**Tip**

You do not want to use the third form when specifying a DEFAULT clause while creating a table. The system will convert now to a timestamp as soon as the constant is parsed, so that when the default value is needed, the time of the table creation would be used! The first two forms will not be evaluated until the default value is used, because they are function calls. Thus they will give the desired behavior of defaulting to the time of row insertion.

### 9.9.5. Delaying Execution

The following functions are available to delay execution of the server process:

pg_sleep(seconds)
pg_sleep_for(interval)
pg_sleep_until(timestamp with time zone)

pg_sleep makes the current session's process sleep until seconds seconds have elapsed. seconds is a value of type double precision, so fractional-second delays can be specified. pg_sleep_for is a convenience function for larger sleep times specified as an interval.
pg_sleep_until is a convenience function for when a specific wake-up time is desired. For example:

```sql
SELECT pg_sleep(1.5);
SELECT pg_sleep_for('5 minutes');
SELECT pg_sleep_until('tomorrow 03:00');
```

**Note**

The effective resolution of the sleep interval is platform-specific; 0.01 seconds is a common value. The sleep delay will be at least as long as specified. It might be longer depending on factors such as server load. In particular, `pg_sleep_until` is not guaranteed to wake up exactly at the specified time, but it will not wake up any earlier.

**Warning**

Make sure that your session does not hold more locks than necessary when calling `pg_sleep` or its variants. Otherwise other sessions might have to wait for your sleeping process, slowing down the entire system.

### 9.10. Enum Support Functions

For enum types (described in Section 8.7), there are several functions that allow cleaner programming without hard-coding particular values of an enum type. These are listed in Table 9.32. The examples assume an enum type created as:

```
CREATE TYPE rainbow AS ENUM ('red', 'orange', 'yellow', 'green', 'blue', 'purple');
```

**Table 9.32. Enum Support Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>enum_first(anyenum)</td>
<td>Returns the first value of the input enum type</td>
<td><code>enum_first(nul-l::rainbow)</code></td>
<td>red</td>
</tr>
<tr>
<td>enum_last(anyenum)</td>
<td>Returns the last value of the input enum type</td>
<td><code>enum_last(nul-l::rainbow)</code></td>
<td>purple</td>
</tr>
<tr>
<td>enum_range(anyenum)</td>
<td>Returns all values of the input enum type in an ordered array</td>
<td><code>enum_range(nul-l::rainbow)</code></td>
<td>{red, orange, yellow, green, blue, purple}</td>
</tr>
<tr>
<td>enum_range(anyenum, anyenum)</td>
<td>Returns the range between the two given enum values, as an ordered array. The values must be from the same enum type. If the first parameter is null, the result will start with the first value of the enum type. If the second parameter is null, the result will end with the</td>
<td><code>enum_range('or-ange'::rainbow, 'green'::rain-bow)</code></td>
<td>{orange, yellow, green}</td>
</tr>
</tbody>
</table>

|   |   | `enum_range(NULL, 'green'::rainbow)` | {red, orange, yellow, green} |
|   |   | `enum_range('orange'::rainbow, NULL)` | {orange, yellow, green, blue, purple} |

257
functions disregard the specific value passed to them; they care only about its declared data type. Either null or a specific value of the type can be passed, with the same result. It is more common to apply these functions to a table column or function argument than to a hardwired type name as suggested by the examples.

### 9.11. Geometric Functions and Operators

The geometric types `point`, `box`, `lseg`, `line`, `path`, `polygon`, and `circle` have a large set of native support functions and operators, shown in Table 9.33, Table 9.34, and Table 9.35.

#### Caution

Note that the “same as” operator, `~=` represents the usual notion of equality for the `point`, `box`, `polygon`, and `circle` types. Some of these types also have an `=` operator, but `=` compares for equal *areas* only. The other scalar comparison operators (`<=` and so on) likewise compare areas for these types.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>Translation</td>
<td><code>box '((0,0),(1,1))' + point '(2.0,0)'</code></td>
</tr>
<tr>
<td>-</td>
<td>Translation</td>
<td><code>box '((0,0),(1,1))' - point '(2.0,0)'</code></td>
</tr>
<tr>
<td>*</td>
<td>Scaling/rotation</td>
<td><code>box '((0,0),(1,1))' * point '(2.0,0)'</code></td>
</tr>
<tr>
<td>/</td>
<td>Scaling/rotation</td>
<td><code>box '((0,0),(2,2))' / point '(2.0,0)'</code></td>
</tr>
<tr>
<td>#</td>
<td>Point or box of intersection</td>
<td><code>box '((1,-1),(-1,1))' # box '((1,1),(-2,-2))'</code></td>
</tr>
<tr>
<td>#</td>
<td>Number of points in path or polygon</td>
<td><code># path '((1,0),(0,1),(-1,0))'</code></td>
</tr>
<tr>
<td>@-@</td>
<td>Length or circumference</td>
<td><code>@-@ path '((0,0),(1,0))'</code></td>
</tr>
<tr>
<td>@@</td>
<td>Center</td>
<td><code>@@ circle '((0,0),10)'</code></td>
</tr>
<tr>
<td>##</td>
<td>Closest point to first operand on second operand</td>
<td><code>point '((0,0))' ## lseg '(((2,0),(0,2))'</code></td>
</tr>
<tr>
<td>&lt;-&gt;</td>
<td>Distance between</td>
<td><code>circle '((0,0),1)' &lt;-&gt; circle '((5,0),1)'</code></td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>Overlaps? (One point in common makes this true.)</td>
<td><code>box '(((0,0),(1,1))' &amp;&amp; box '(((0,0),(2,2))'</code></td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>Is strictly left of?</td>
<td><code>circle '((0,0),1)' &lt;&lt; circle '((5,0),1)'</code></td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>Is strictly right of?</td>
<td><code>circle '((5,0),1)' &gt;&gt; circle '((0,0),1)'</code></td>
</tr>
<tr>
<td>Operator</td>
<td>Description</td>
<td>Example</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td>Does not extend to the right of?</td>
<td>box '((0,0),(1,1))' &amp;&lt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&amp;&gt;</td>
<td>Does not extend to the left of?</td>
<td>box '((0,0),(3,3))' &amp;&gt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>Is strictly below?</td>
<td>box '((0,0),(3,3))' &lt;&lt; box '((3,4),(5,5))'</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>Is strictly above?</td>
<td>box '((3,4),(5,5))' &gt;&gt; box '((0,0),(3,3))'</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td>Does not extend above?</td>
<td>box '((0,0),(1,1))' &amp;&lt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&amp;&gt;</td>
<td>Does not extend below?</td>
<td>box '((0,0),(3,3))' &amp;&gt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>^&lt;</td>
<td>Is below (allows touching)?</td>
<td>circle '((0,0),1)' ^&lt; circle '((0,5),1)'</td>
</tr>
<tr>
<td>^&gt;</td>
<td>Is above (allows touching)?</td>
<td>circle '((0,5),1)' ^&gt; circle '((0,0),1)'</td>
</tr>
<tr>
<td>#?</td>
<td>Intersect?</td>
<td>lseg '((-1,0),(1,0))' #? box '((-2,-2),(2,2))'</td>
</tr>
<tr>
<td>?-</td>
<td>Is horizontal?</td>
<td>?- lseg '((-1,0),(1,0))'</td>
</tr>
<tr>
<td>?-</td>
<td>Are horizontally aligned?</td>
<td>point '(1,0)' ?- point '(0,0)'</td>
</tr>
<tr>
<td>?</td>
<td></td>
<td>Is vertical?</td>
</tr>
<tr>
<td>?</td>
<td></td>
<td>Are vertically aligned?</td>
</tr>
<tr>
<td>?-</td>
<td></td>
<td>Is perpendicular?</td>
</tr>
<tr>
<td>?</td>
<td></td>
<td>Are parallel?</td>
</tr>
<tr>
<td>@&gt;</td>
<td>Contains?</td>
<td>circle '((0,0),2)' @&gt; point '(1,1)'</td>
</tr>
<tr>
<td>&lt;@</td>
<td>Contained in or on?</td>
<td>point '(1,1)' &lt;@ circle '((0,0),2)'</td>
</tr>
<tr>
<td>-=</td>
<td>Same as?</td>
<td>polygon '((0,0),(1,1))' -= polygon '((1,1),(0,0))'</td>
</tr>
</tbody>
</table>

**Note**

Before PostgreSQL 8.2, the containment operators @> and <@ were respectively called ~ and @. These names are still available, but are deprecated and will eventually be removed.
### Table 9.34. Geometric Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>area(object)</td>
<td>double precision</td>
<td>area</td>
<td>area(box '((0,0),(1,1))')</td>
</tr>
<tr>
<td>center(object)</td>
<td>point</td>
<td>center</td>
<td>center(box '((0,0),(1,2))')</td>
</tr>
<tr>
<td>diameter(circle)</td>
<td>double precision</td>
<td>diameter of circle</td>
<td>diameter(circle '((0,0),2.0)')</td>
</tr>
<tr>
<td>height(box)</td>
<td>double precision</td>
<td>vertical size of box</td>
<td>height(box '((0,0),(1,1))')</td>
</tr>
<tr>
<td>isclosed(path)</td>
<td>boolean</td>
<td>a closed path?</td>
<td>isclosed(path '((0,0),(1,1), (2,0))')</td>
</tr>
<tr>
<td>isopen(path)</td>
<td>boolean</td>
<td>an open path?</td>
<td>isopen(path '[(0,0),(1,1), (2,0)]')</td>
</tr>
<tr>
<td>length(object)</td>
<td>double precision</td>
<td>length</td>
<td>length(path '((-1,0), (1,0))')</td>
</tr>
<tr>
<td>npoints(path)</td>
<td>int</td>
<td>number of points</td>
<td>npoints(path '[(0,0),(1,1), (2,0)])')</td>
</tr>
<tr>
<td>npoints(polygon)</td>
<td>int</td>
<td>number of points</td>
<td>npoints(polygon '((1,1), (0,0))')</td>
</tr>
<tr>
<td>pclose(path)</td>
<td>path</td>
<td>convert path to closed</td>
<td>pclose(path '[(0,0),(1,1), (2,0)])')</td>
</tr>
<tr>
<td>open(path)</td>
<td>path</td>
<td>convert path to open</td>
<td>open(path '[(0,0),(1,1), (2,0)])')</td>
</tr>
<tr>
<td>radius(circle)</td>
<td>double precision</td>
<td>radius of circle</td>
<td>radius(circle '((0,0),2.0)')</td>
</tr>
<tr>
<td>width(box)</td>
<td>double precision</td>
<td>horizontal size of box</td>
<td>width(box '((0,0),(1,1))')</td>
</tr>
</tbody>
</table>

### Table 9.35. Geometric Type Conversion Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>box(circle)</td>
<td>box</td>
<td>circle to box</td>
<td>box(circle '((0,0),2.0)')</td>
</tr>
<tr>
<td>box(point)</td>
<td>box</td>
<td>point to empty box</td>
<td>box(point '(0,0)')</td>
</tr>
<tr>
<td>box(point, point)</td>
<td>box</td>
<td>points to box</td>
<td>box(point '(0,0)', point '(1,1)')</td>
</tr>
<tr>
<td>box(polygon)</td>
<td>box</td>
<td>polygon to box</td>
<td>box(polygon '((0,0),(1,1), (2,0))')</td>
</tr>
<tr>
<td>bound_box(box, box)</td>
<td>box</td>
<td>boxes to bounding box</td>
<td>bound_box(box '((0,0),(1,1))')</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
</tr>
<tr>
<td>--------------------------------</td>
<td>-------------</td>
<td>------------------------</td>
<td>----------------------------------</td>
</tr>
<tr>
<td>circle(box)</td>
<td>circle</td>
<td>box to circle</td>
<td>circle(box '((3,3), (4,4))')</td>
</tr>
<tr>
<td>circle(point, double precision)</td>
<td>circle</td>
<td>center and radius to circle</td>
<td>circle(point '(-1,0)', 2.0)</td>
</tr>
<tr>
<td>circle(polygone)</td>
<td>circle</td>
<td>polygon to circle</td>
<td>circle(polygone '((-1,0), (1,1), (2,0))')</td>
</tr>
<tr>
<td>line(point, point)</td>
<td>line</td>
<td>points to line</td>
<td>line(point '(-1,0)', point '(-1,0)')</td>
</tr>
<tr>
<td>lseg(box)</td>
<td>lseg</td>
<td>box diagonal to line segment</td>
<td>lseg(box '((-1,0), (1,0))')</td>
</tr>
<tr>
<td>lseg(point, point)</td>
<td>lseg</td>
<td>points to line segment</td>
<td>lseg(point '(-1,0)', point '(-1,0)')</td>
</tr>
<tr>
<td>path(polygone)</td>
<td>path</td>
<td>polygon to path</td>
<td>path(polygone '((-1,0), (1,1), (2,0))')</td>
</tr>
<tr>
<td>point(double precision, double precision)</td>
<td>point</td>
<td>construct point</td>
<td>point(23.4, -44.5)</td>
</tr>
<tr>
<td>point(box)</td>
<td>point</td>
<td>center of box</td>
<td>point(box '((-1,0), (1,0))')</td>
</tr>
<tr>
<td>point(circle)</td>
<td>point</td>
<td>center of circle</td>
<td>point(circle '((-1,0), (1,0))')</td>
</tr>
<tr>
<td>point(lseg)</td>
<td>point</td>
<td>center of line segment</td>
<td>point(lseg '((-1,0), (1,0))')</td>
</tr>
<tr>
<td>point(polygone)</td>
<td>point</td>
<td>center of polygon</td>
<td>point(polygone '((-1,0), (1,1), (2,0))')</td>
</tr>
<tr>
<td>polygon(box)</td>
<td>polygon</td>
<td>box to 4-point polygon</td>
<td>polygon(box '((-1,0), (1,1))')</td>
</tr>
<tr>
<td>polygon(circle)</td>
<td>polygon</td>
<td>circle to 12-point polygon</td>
<td>polygon(circle '((-1,0), (1,1))')</td>
</tr>
<tr>
<td>polygon(npts, circle)</td>
<td>polygon</td>
<td>circle to npts-point polygon</td>
<td>polygon(12, circle '((-1,0), (1,1))')</td>
</tr>
<tr>
<td>polygon(path)</td>
<td>polygon</td>
<td>path to polygon</td>
<td>polygon(path '((-1,0), (1,1), (2,0))')</td>
</tr>
</tbody>
</table>

It is possible to access the two component numbers of a point as though the point were an array with indexes 0 and 1. For example, if t.p is a point column then SELECT p[0] FROM t retrieves...
the X coordinate and UPDATE t SET p[1] = ... changes the Y coordinate. In the same way, a value of type box or lseg can be treated as an array of two point values.

The area function works for the types box, circle, and path. The area function only works on the path data type if the points in the path are non-intersecting. For example, the path '((0,0), (0,1), (2,1), (2,2), (1,2), (1,0), (0,0))'::PATH will not work; however, the following visually identical path '((0,0), (0,1), (1,1), (1,2), (2,2), (2,1), (1,1), (1,0), (0,0))'::PATH will work. If the concept of an intersecting versus non-intersecting path is confusing, draw both of the above paths side by side on a piece of graph paper.

9.12. Network Address Functions and Operators

Table 9.36 shows the operators available for the cidr and inet types. The operators <<, <<=, >>, and && test for subnet inclusion. They consider only the network parts of the two addresses (ignoring any host part) and determine whether one network is identical to or a subnet of the other.

### Table 9.36. cidr and inet Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;</td>
<td>is less than</td>
<td>inet '192.168.1.5' &lt; inet '192.168.1.6'</td>
</tr>
<tr>
<td>&lt;=</td>
<td>is less than or equal</td>
<td>inet '192.168.1.5' &lt;= inet '192.168.1.5'</td>
</tr>
<tr>
<td>=</td>
<td>equals</td>
<td>inet '192.168.1.5' = inet '192.168.1.5'</td>
</tr>
<tr>
<td>&gt;=</td>
<td>is greater or equal</td>
<td>inet '192.168.1.5' &gt;= inet '192.168.1.5'</td>
</tr>
<tr>
<td>&gt;</td>
<td>is greater than</td>
<td>inet '192.168.1.5' &gt; inet '192.168.1.4'</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>is not equal</td>
<td>inet '192.168.1.5' &lt;&gt; inet '192.168.1.4'</td>
</tr>
<tr>
<td>&lt;&lt;=</td>
<td>is contained by</td>
<td>inet '192.168.1.5' &lt;&lt;= inet '192.168.1/24'</td>
</tr>
<tr>
<td>&lt;=</td>
<td>is contained by or equals</td>
<td>inet '192.168.1/24' &lt;= inet '192.168.1/24'</td>
</tr>
<tr>
<td>&gt;&gt;&gt;</td>
<td>contains</td>
<td>inet '192.168.1/24' &gt;&gt;&gt; inet '192.168.1.5'</td>
</tr>
<tr>
<td>&gt;=</td>
<td>contains or equals</td>
<td>inet '192.168.1/24' &gt;= inet '192.168.1.5'</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>contains or is contained by</td>
<td>inet '192.168.1/24' &amp;&amp; inet '192.168.1/24'</td>
</tr>
<tr>
<td>~</td>
<td>bitwise NOT</td>
<td>~ inet '192.168.1.6'</td>
</tr>
<tr>
<td>&amp;</td>
<td>bitwise AND</td>
<td>inet '192.168.1.6' &amp; inet '0.0.0.255'</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bitwise OR</td>
</tr>
</tbody>
</table>
## Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>addition</td>
<td>inet '192.168.1.6' + 25</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>inet '192.168.1.43' - 36</td>
</tr>
<tr>
<td></td>
<td>subtraction</td>
<td>inet '192.168.1.43' - inet '192.168.1.19'</td>
</tr>
</tbody>
</table>

Table 9.37 shows the functions available for use with the `cidr` and `inet` types. The `abbrev`, `host`, and `text` functions are primarily intended to offer alternative display formats.

### Table 9.37. `cidr` and `inet` Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>abbrev(inet)</code></td>
<td><code>text</code></td>
<td>abbreviated display format as text</td>
<td><code>abbrev(inet '10.1.0.0/16')</code></td>
<td><code>10.1.0.0/16</code></td>
</tr>
<tr>
<td><code>abbrev(cidr)</code></td>
<td><code>text</code></td>
<td>abbreviated display format as text</td>
<td><code>abbrev(cidr '10.1.16')</code></td>
<td><code>10.1/16</code></td>
</tr>
<tr>
<td><code>broadcast(inet)</code></td>
<td><code>inet</code></td>
<td>broadcast address for network</td>
<td><code>broadcast('192.168.1.5/24')</code></td>
<td><code>192.168.1.255/24</code></td>
</tr>
<tr>
<td><code>family(inet)</code></td>
<td><code>int</code></td>
<td>extract family of address; 4 for IPv4, 6 for IPv6</td>
<td><code>family('::1')</code></td>
<td>6</td>
</tr>
<tr>
<td><code>host(inet)</code></td>
<td><code>text</code></td>
<td>extract IP address as text</td>
<td><code>host('192.168.1.5/24')</code></td>
<td><code>192.168.1.5</code></td>
</tr>
<tr>
<td><code>hostmask(inet)</code></td>
<td><code>inet</code></td>
<td>construct host mask for network</td>
<td><code>hostmask('192.168.23.20/30')</code></td>
<td><code>0.0.0.3</code></td>
</tr>
<tr>
<td><code>masklen(inet)</code></td>
<td><code>int</code></td>
<td>extract netmask length</td>
<td><code>masklen('192.168.1.5/24')</code></td>
<td>24</td>
</tr>
<tr>
<td><code>netmask(inet)</code></td>
<td><code>inet</code></td>
<td>construct netmask for network</td>
<td><code>netmask('192.168.1.5/24')</code></td>
<td><code>255.255.255.0</code></td>
</tr>
<tr>
<td><code>network(in-cidr)</code></td>
<td><code>inet</code></td>
<td>extract network part of address</td>
<td><code>network('192.168.1.5/24')</code></td>
<td><code>192.168.1.0/24</code></td>
</tr>
<tr>
<td><code>set_masklen(inet, int)</code></td>
<td><code>inet</code></td>
<td>set netmask length for inet value</td>
<td><code>set_masklen('192.168.1.5/24', 16)</code></td>
<td><code>192.168.0.0/16</code></td>
</tr>
<tr>
<td><code>set_masklen(cidr, int)</code></td>
<td><code>inet</code></td>
<td>set netmask length for cidr value</td>
<td><code>set_masklen('192.168.1.5/24', 16)</code></td>
<td><code>192.168.1.0/24</code></td>
</tr>
<tr>
<td><code>text(inet)</code></td>
<td><code>inet</code></td>
<td>extract IP address and netmask length as text</td>
<td><code>text(inet '192.168.1.5')</code></td>
<td><code>192.168.1.5/32</code></td>
</tr>
<tr>
<td><code>in boolean</code></td>
<td><code>boolean</code></td>
<td>are the addresses in the same family?</td>
<td><code>in(boolean '192.168.1.5/24', '192.168.1.5/24')</code></td>
<td><code>false</code></td>
</tr>
<tr>
<td><code>inet_same_family(inet, inet)</code></td>
<td><code>boolean</code></td>
<td>are the addresses in the same family?</td>
<td><code>inet_same_family('192.168.1.5/24', '192.168.1.5/24')</code></td>
<td><code>false</code></td>
</tr>
<tr>
<td><code>inet_merge(inet, inet)</code></td>
<td><code>inet</code></td>
<td>the smallest network which includes both of the given networks</td>
<td><code>inet_merge('192.168.1.5/24', '192.168.0.0/16')</code></td>
<td><code>192.168.0.0/22</code></td>
</tr>
</tbody>
</table>
Any `cidr` value can be cast to `inet` implicitly or explicitly; therefore, the functions shown above as operating on `inet` also work on `cidr` values. (Where there are separate functions for `inet` and `cidr`, it is because the behavior should be different for the two cases.) Also, it is permitted to cast an `inet` value to `cidr`. When this is done, any bits to the right of the netmask are silently zeroed to create a valid `cidr` value. In addition, you can cast a text value to `inet` or `cidr` using normal casting syntax: for example, `inet(expression)` or `colname::cidr`.

Table 9.38 shows the functions available for use with the `macaddr` type. The function `trunc(macaddr)` returns a MAC address with the last 3 bytes set to zero. This can be used to associate the remaining prefix with a manufacturer.

### Table 9.38. `macaddr` Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>trunc(macaddr)</code></td>
<td><code>macaddr</code></td>
<td>set last 3 bytes to zero</td>
<td><code>trunc(macaddr '12:34:56:78:90:ab')</code></td>
<td><code>12:34:56:00:00:00</code></td>
</tr>
</tbody>
</table>

The `macaddr` type also supports the standard relational operators (>, <=, etc.) for lexicographical ordering, and the bitwise arithmetic operators (~, &, and |) for NOT, AND and OR.

Table 9.39 shows the functions available for use with the `macaddr8` type. The function `trunc(macaddr8)` returns a MAC address with the last 5 bytes set to zero. This can be used to associate the remaining prefix with a manufacturer.

### Table 9.39. `macaddr8` Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>trunc(macaddr8)</code></td>
<td><code>macaddr8</code></td>
<td>set last 5 bytes to zero</td>
<td><code>trunc(macaddr8 '12:34:56:78:90:ab:cd:ef')</code></td>
<td><code>12:34:56:00:00:00:00:00</code></td>
</tr>
<tr>
<td><code>macaddr8_set7bit(macaddr8)</code></td>
<td><code>macaddr8</code></td>
<td>set 7th bit to one, also known as modified EUI-64, for inclusion in an IPv6 address</td>
<td><code>macaddr8_set7bit(macaddr8 '00:34:56:ab:cd:ef')</code></td>
<td><code>02:34:56:ff:fe:ab:cd:ef</code></td>
</tr>
</tbody>
</table>

The `macaddr8` type also supports the standard relational operators (>, <=, etc.) for ordering, and the bitwise arithmetic operators (~, &, and |) for NOT, AND and OR.

### 9.13. Text Search Functions and Operators

Table 9.40, Table 9.41 and Table 9.42 summarize the functions and operators that are provided for full text searching. See Chapter 12 for a detailed explanation of PostgreSQL’s text search facility.

### Table 9.40. Text Search Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>@@</td>
<td>boolean</td>
<td><code>tsvector</code> matches query?</td>
<td><code>to_tsvector('fat cats ate rats')</code></td>
<td>t</td>
</tr>
<tr>
<td>@@@</td>
<td>boolean</td>
<td>deprecated synonym for @@</td>
<td><code>to_tsvector('fat cats')</code></td>
<td>t</td>
</tr>
</tbody>
</table>
### Functions and Operators

#### Table 9.41. Text Search Functions

<table>
<thead>
<tr>
<th>Operator</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>`</td>
<td></td>
<td>`</td>
<td>tsvector</td>
<td>concatenate tsvectors</td>
</tr>
<tr>
<td><code>&amp;&amp;</code></td>
<td>tsquery</td>
<td>AND tsqueries together</td>
<td>'fat':tsquery</td>
<td></td>
</tr>
<tr>
<td>`</td>
<td></td>
<td>`</td>
<td>tsquery</td>
<td>OR tsqueries together</td>
</tr>
<tr>
<td><code>!!</code></td>
<td>tsquery</td>
<td>negate tsquery</td>
<td>!! 'cat':tsquery</td>
<td></td>
</tr>
<tr>
<td><code>&lt;-&gt;</code></td>
<td>tsquery</td>
<td>tsquery followed by tsquery</td>
<td>to_tsquery('fat') 'rat'</td>
<td></td>
</tr>
<tr>
<td><code>@&gt;</code></td>
<td>boolean</td>
<td>tsquery contains another?</td>
<td>'cat':tsquery</td>
<td></td>
</tr>
<tr>
<td><code>&lt;@</code></td>
<td>boolean</td>
<td>tsquery is contained in?</td>
<td>'cat':tsquery</td>
<td></td>
</tr>
</tbody>
</table>

### Note

The `tsquery` containment operators consider only the lexemes listed in the two queries, ignoring the combining operators.

In addition to the operators shown in the table, the ordinary B-tree comparison operators (=, <, etc) are defined for types `tsvector` and `tsquery`. These are not very useful for text searching but allow, for example, unique indexes to be built on columns of these types.

---

**Table 9.41. Text Search Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>array_to_tsvector(text[])</code></td>
<td>tsvector</td>
<td>convert array of lexemes to tsvector</td>
<td><code>array((fat, cat, rat))::text[]</code></td>
<td><code>fat' 'cat' 'rat'</code></td>
</tr>
</tbody>
</table>

---

Page 265
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>get_current_ts_configuration()</td>
<td>regconfig</td>
<td>get default text search configuration</td>
<td>get_current_ts_configuration()</td>
<td>english</td>
</tr>
<tr>
<td>length(tsvector)</td>
<td>integer</td>
<td>number of lexemes</td>
<td>length('fat:2,4 cat:3 rat:5A':tsvector)</td>
<td></td>
</tr>
<tr>
<td>numnode(ts_query)</td>
<td>integer</td>
<td>number of lexemes plus operators in tsquery</td>
<td>numnode('(fat &amp; rat)</td>
<td>cat':tsquery)</td>
</tr>
<tr>
<td>plainto_tsquery(query({config regconfig, } query text))</td>
<td>tsquery</td>
<td>produce tsquery ignoring punctuation</td>
<td>plainto_tsquery('english', 'The Fat Rats')</td>
<td></td>
</tr>
<tr>
<td>phraseto_tsquery(query({config regconfig, } query text))</td>
<td>tsquery</td>
<td>produce tsquery that searches for a phrase, ignoring punctuation</td>
<td>phraseto_tsquery('english', 'The Fat Rats')</td>
<td></td>
</tr>
<tr>
<td>websearch_to_tsquery(query({config regconfig, } query text))</td>
<td>tsquery</td>
<td>produce tsquery from a web search style query</td>
<td>websearch_to_tsquery('english', &quot;fat rat&quot;</td>
<td>'rat')</td>
</tr>
<tr>
<td>querytree(tsquery)</td>
<td>text</td>
<td>get indexable part of a tsquery tree('foo &amp; ! bar':tsquery)</td>
<td>querytree('foo &amp; ! bar':tsquery)</td>
<td></td>
</tr>
<tr>
<td>setweight(vector, weight &quot;char&quot;)</td>
<td>tsvector</td>
<td>assign weight to each element of vector</td>
<td>setweight('fat:2,4A cat:3 rat:5B':tsvector:5A 'tor, 'A')</td>
<td></td>
</tr>
<tr>
<td>setweight(vector, weight &quot;char&quot;, lexemes text[])</td>
<td>tsvector</td>
<td>assign weight to elements of vector that are listed in lexemes</td>
<td>setweight('fat:2,4 cat:3 rat:5B':tsvector:5A 'tor, 'A', '(cat, rat)')</td>
<td></td>
</tr>
<tr>
<td>strip(tsvector)</td>
<td>tsvector</td>
<td>remove positions and weights from tsvector</td>
<td>strip('fat:2,4 cat:3 rat:5B':tsvector:5A 'tor, 'A')</td>
<td></td>
</tr>
<tr>
<td>to_tsquery([tsquery config regconfig, } query text)</td>
<td>tsvector</td>
<td>normalize words and convert to tsquery</td>
<td>to_tsquery('english', 'The &amp; Fat &amp; Rats')</td>
<td></td>
</tr>
<tr>
<td>to_tsvector(vector)</td>
<td>tsvector</td>
<td>reduce document to tsvector text to tsvector</td>
<td>to_tsvector('fat:2')</td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------------------------------</td>
<td>-------------</td>
<td>-------------</td>
<td>--------------------------</td>
<td>-------------------</td>
</tr>
<tr>
<td><code>fig regconfig document text</code></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>to_tsvector([  config  ] document json(b))</code></td>
<td></td>
<td>reduce each string value in the document, then concatenate those in document order to produce a single tsvector</td>
<td><code>to_tsvector('english', '{&quot;a&quot;: &quot;The Fat Rats&quot;}')</code></td>
<td><code>{&quot;fat&quot;:2, &quot;rat&quot;:3}</code></td>
</tr>
<tr>
<td><code>json(b)_to_tsvector([  config  ] document json(b), filter json(b))</code></td>
<td></td>
<td>reduce each value in the document, specified by filter, then concatenate those in document order to produce a single tsvector</td>
<td><code>json_to_tsvector('english', '{&quot;a&quot;: &quot;The Fat Rats&quot;, &quot;b&quot;: 123}':json, 'string', 'numeric')</code></td>
<td><code>{&quot;fat&quot;:2, &quot;rat&quot;:3, &quot;123&quot;:5}</code></td>
</tr>
<tr>
<td><code>ts_delete(vector tsvector, lexeme text)</code></td>
<td><code>vector</code></td>
<td>remove given lexeme from vector</td>
<td><code>ts_delete('fat:2,4 cat:3 rat:5A':tsvector, 'fat')</code></td>
<td><code>{&quot;cat&quot;:3, &quot;rat&quot;:5A}</code></td>
</tr>
<tr>
<td><code>ts_delete(vector tsvector)</code></td>
<td></td>
<td>remove any occurrence of lexemes</td>
<td><code>ts_delete('fat:2,4 cat:3 rat:5A':tsvector)</code></td>
<td><code>{&quot;cat&quot;:3}</code></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------------------</td>
<td>-------------</td>
<td>--------------------------------------</td>
<td>--------------------------------</td>
<td>-----------------</td>
</tr>
<tr>
<td><code>tor, lexemes</code></td>
<td></td>
<td>in <code>lexemes</code> from <code>rat:5A</code>::tsvecctor, AR-RAY['fat','rat']</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_fil</code></td>
<td>tsvector</td>
<td>select only elements with given weights from <code>rat:3B</code>::tsvector, AR-RAY['fat','rat']</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_head</code></td>
<td>text</td>
<td>display a query match</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_head</code></td>
<td>text</td>
<td>display a query match</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_rank</code></td>
<td>float4</td>
<td>rank document for query</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_rank_cd</code></td>
<td>float4</td>
<td>rank document for query using cover density</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_rewrite</code></td>
<td>tsquery</td>
<td>replace target with substitute within query</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ts_rewrite</code></td>
<td>tsquery</td>
<td>replace using target from a SELECT &amp;</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

268
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ts-</code> <code>tsquery</code> query_phrase(query1 <code>tsquery</code>, query2 <code>tsquery</code>)</td>
<td><code>tsquery</code></td>
<td>make query that searches for <code>query1</code> followed by <code>query2</code> (same as <code>&lt;&gt;</code> operator)</td>
<td><code>to_tsquery('fat')</code></td>
<td><code>&lt;&gt;</code></td>
</tr>
<tr>
<td><code>ts-</code> <code>tsquery</code> query_phrase(query1 <code>tsquery</code>, query2 <code>tsquery</code>, distance integer)</td>
<td><code>tsquery</code></td>
<td>make query that searches for <code>query1</code> followed by <code>query2</code> at distance <code>distance</code></td>
<td><code>to_tsquery('fat')</code></td>
<td><code>&lt;10&gt;</code></td>
</tr>
<tr>
<td><code>tsvector_to_array</code> (tsvector)</td>
<td>text[]</td>
<td>convert tsvector to array of lexemes</td>
<td><code>to_tsquery('fat:2,4 cat:3 rat:5A'::tsvector)</code></td>
<td><code>{cat,fat,rat}</code></td>
</tr>
<tr>
<td><code>tsvector_update_trigger</code> date_trig- ger()</td>
<td>trigger function for automatic tsvector column update</td>
<td><code>CREATE TRIGGER</code> ... <code>tsvector_update_trig- ger(tsvcol, 'pg_cata- log.swedish', title, body)</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>tsvector_update_trigger</code> date_trig- ger_column()</td>
<td>trigger function for automatic tsvector column update</td>
<td><code>CREATE TRIGGER</code> ... <code>tsvector_update_trig- ger_column(tsvcol, configcol, title, body)</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>unnest</code> (tsvector, OUT lexeme text, OUT positions small-int[], OUT weights text)</td>
<td>setof record</td>
<td>expand a tsvector to a set of rows</td>
<td><code>unnest('fat:2,4 cat:3, {D})</code></td>
<td></td>
</tr>
</tbody>
</table>

**Note**

All the text search functions that accept an optional `regconfig` argument will use the configuration specified by `default_text_search_config` when that argument is omitted.
The functions in Table 9.42 are listed separately because they are not usually used in everyday text searching operations. They are helpful for development and debugging of new text search configurations.

### Table 9.42. Text Search Debugging Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ts_debug()</code></td>
<td><code>setof record</code></td>
<td>test a configuration</td>
<td><code>ts_debug('english', 'The brightest supernovae')</code></td>
<td>ascii-word,&quot;Word, English&quot;, 'The all brightest supernova'</td>
</tr>
<tr>
<td><code>ts_lexize()</code></td>
<td><code>text[]</code></td>
<td>test a dictionary</td>
<td><code>ts_lexize('english_stem', 'stars')</code></td>
<td><code>{'star'}</code></td>
</tr>
<tr>
<td><code>ts_parse()</code></td>
<td><code>setof record</code></td>
<td>test a parser</td>
<td><code>ts_parse('default', 'foo - bar')</code></td>
<td><code>(1,foo)</code></td>
</tr>
<tr>
<td><code>ts_parse()</code></td>
<td><code>setof record</code></td>
<td>test a parser</td>
<td><code>ts_parse(3722, 'foo - bar')</code></td>
<td><code>(1,foo)</code></td>
</tr>
<tr>
<td><code>ts_tooken_type()</code></td>
<td><code>setof record</code></td>
<td>get token types defined by parser</td>
<td><code>ts_tooken_type('default')</code></td>
<td><code>(1,ascii-word,&quot;Word,&quot; default all ASCII&quot;)</code></td>
</tr>
<tr>
<td><code>ts_tooken_type()</code></td>
<td><code>setof record</code></td>
<td>get token types defined by parser</td>
<td><code>ts_tooken_type(3722)</code></td>
<td><code>(1,ascii-word,&quot;Word,&quot; 3722 all ASCII&quot;)</code></td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>OUT tokid in-</td>
<td>integer, OUT alias text, OUT description text</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ts_stat(sql-setof record query text, [weights text,] OUT word text, OUT ndoc integer, OUT nentry integer)</td>
<td></td>
<td>get statistics of ts_stat('S- (foo,10,15) ... a tsvector column vector unm from apod')</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

### 9.14. XML Functions

The functions and function-like expressions described in this section operate on values of type xml. Check Section 8.13 for information about the xml type. The function-like expressions xmlparse and xmlserialize for converting to and from type xml are not repeated here. Use of most of these functions requires the installation to have been built with configure --with-libxml.

#### 9.14.1. Producing XML Content

A set of functions and function-like expressions are available for producing XML content from SQL data. As such, they are particularly suitable for formatting query results into XML documents for processing in client applications.

##### 9.14.1.1. xmlcomment

xmlcomment (text)

The function xmlcomment creates an XML value containing an XML comment with the specified text as content. The text cannot contain "--" or end with a "-" so that the resulting construct is a valid XML comment. If the argument is null, the result is null.

Example:

```
SELECT xmlcomment('hello');
```

xmlcomment

```
<!--hello-->
```

##### 9.14.1.2. xmlconcat

xmlconcat (xml[, ...])

The function xmlconcat concatenates a list of individual XML values to create a single value containing an XML content fragment. Null values are omitted; the result is only null if there are no non-null arguments.

Example:

```sql
xmlconcat(xml[,'...']);
```
xmlconcat
----------------------
<abc/> <bar>foo</bar>

XML declarations, if present, are combined as follows. If all argument values have the same XML version declaration, that version is used in the result, else no version is used. If all argument values have the standalone declaration value "yes", then that value is used in the result. If all argument values have a standalone declaration value and at least one is "no", then that is used in the result. Else the result will have no standalone declaration. If the result is determined to require a standalone declaration but no version declaration, a version declaration with version 1.0 will be used because XML requires an XML declaration to contain a version declaration. Encoding declarations are ignored and removed in all cases.

Example:

xmlconcat
----------------------
<?xml version="1.1"?><foo/> <bar/>

9.14.1.3. xmlelement

xmlelement (name name [, xmlattributes (value [AS attname] [, ... ])] [, content, ...])

The xmlelement expression produces an XML element with the given name, attributes, and content.

Examples:

SELECT xmlelement (name foo);

SELECT xmlelement (name foo, xmlattributes ('xyz' as bar));

SELECT xmlelement (name foo, xmlattributes (current_date as bar), 'cont', 'ent');

Element and attribute names that are not valid XML names are escaped by replacing the offending characters by the sequence _xHHHH_, where HHHH is the character's Unicode codepoint in hexadecimal notation. For example:
SELECT xmlelement(name "foo$bar", xmlattributes('xyz' as "a&b"));

---
<foo_x0024_bar a_x0026_b="xyz"/>

An explicit attribute name need not be specified if the attribute value is a column reference, in which case the column's name will be used as the attribute name by default. In other cases, the attribute must be given an explicit name. So this example is valid:

CREATE TABLE test (a xml, b xml);
SELECT xmlelement(name test, xmlattributes(a, b)) FROM test;

But these are not:

SELECT xmlelement(name test, xmlattributes('constant'), a, b) FROM test;
SELECT xmlelement(name test, xmlattributes(func(a, b))) FROM test;

Element content, if specified, will be formatted according to its data type. If the content is itself of type xml, complex XML documents can be constructed. For example:

SELECT xmlelement(name foo, xmlattributes('xyz' as bar),
                     xmlelement(name abc),
                     xmlcomment('test'),
                     xmlelement(name xyz));

---
<foo bar="xyz"><abc/><!--test--><xyz/></foo>

Content of other types will be formatted into valid XML character data. This means in particular that the characters <, >, and & will be converted to entities. Binary data (data type bytea) will be represented in base64 or hex encoding, depending on the setting of the configuration parameter xmlbinary. The particular behavior for individual data types is expected to evolve in order to align the SQL and PostgreSQL data types with the XML Schema specification, at which point a more precise description will appear.

9.14.1.4. xmlforest

xmlforest (content [AS name] [, ...])

The xmlforest expression produces an XML forest (sequence) of elements using the given names and content.

Examples:

SELECT xmlforest('abc' AS foo, 123 AS bar);

---
<foo>abc</foo><bar>123</bar>
Functions and Operators

SELECT xmlforest(table_name, column_name)
FROM information_schema.columns
WHERE table_schema = 'pg_catalog';

xmlforest

sciousName>pg_authid</table_name><column_name>rolname</column_name>
sciousName>pg_authid</table_name><column_name>rolsuper</column_name>
...

As seen in the second example, the element name can be omitted if the content value is a column reference, in which case the column name is used by default. Otherwise, a name must be specified.

Element names that are not valid XML names are escaped as shown for xmlelement above. Similarly, content data is escaped to make valid XML content, unless it is already of type xml.

Note that XML forests are not valid XML documents if they consist of more than one element, so it might be useful to wrap xmlforest expressions in xmlelement.

9.14.1.5. xmlpi

xmlpi(name target [, content])

The xmlpi expression creates an XML processing instruction. The content, if present, must not contain the character sequence `?>`.

Example:

SELECT xmlpi(name php, 'echo "hello world";');

xmlpi

----------
<?php echo "hello world";?><

9.14.1.6. xmlroot

xmlroot(xml, version text | no value [, standalone yes|no|no
value])

The xmlroot expression alters the properties of the root node of an XML value. If a version is specified, it replaces the value in the root node’s version declaration; if a standalone setting is specified, it replaces the value in the root node’s standalone declaration.

SELECT xmlroot(xmlparse(document '<?xml version="1.1"?
><content>abc</content>'),
   version '1.0', standalone yes);

xmlroot

--------------------
<?xml version="1.0" standalone="yes"?><
<content>abc</content>
9.14.1.7. \texttt{xmlagg}

\texttt{xmlagg(xml)}

The function \texttt{xmlagg} is, unlike the other functions described here, an aggregate function. It concatenates the input values to the aggregate function call, much like \texttt{xmlconcat} does, except that concatenation occurs across rows rather than across expressions in a single row. See Section 9.20 for additional information about aggregate functions.

Example:

\begin{verbatim}
CREATE TABLE test (y int, x xml);
INSERT INTO test VALUES (1, '<foo>abc</foo>');
INSERT INTO test VALUES (2, '<bar/>');
SELECT xmlagg(x) FROM test;
\end{verbatim}

\begin{verbatim}
----------------------
<foo>abc</foo><bar/>
\end{verbatim}

To determine the order of the concatenation, an \texttt{ORDER BY} clause may be added to the aggregate call as described in Section 4.2.7. For example:

\begin{verbatim}
SELECT xmlagg(x ORDER BY y DESC) FROM test;
\end{verbatim}

\begin{verbatim}
----------------------
<bar/> <foo>abc</foo>
\end{verbatim}

The following non-standard approach used to be recommended in previous versions, and may still be useful in specific cases:

\begin{verbatim}
SELECT xmlagg(x) FROM (SELECT * FROM test ORDER BY y DESC) AS tab;
\end{verbatim}

\begin{verbatim}
----------------------
<bar/> <foo>abc</foo>
\end{verbatim}

9.14.2. XML Predicates

The expressions described in this section check properties of \texttt{xml} values.

9.14.2.1. \texttt{IS DOCUMENT}

\texttt{xml IS DOCUMENT}

The expression \texttt{IS DOCUMENT} returns true if the argument XML value is a proper XML document, false if it is not (that is, it is a content fragment), or null if the argument is null. See Section 8.13 about the difference between documents and content fragments.

9.14.2.2. \texttt{IS NOT DOCUMENT}

\texttt{xml IS NOT DOCUMENT}

The expression \texttt{IS NOT DOCUMENT} returns false if the argument XML value is a proper XML document, true if it is not (that is, it is a content fragment), or null if the argument is null.
### 9.14.2.3. `xmlexists`

The function `xmlexists` returns true if the XPath expression in the first argument returns any nodes, and false otherwise. (If either argument is null, the result is null.)

Example:

```sql
SELECT xmlexists('//town[text() = 'Toronto']' PASSING BY REF '<towns><town>Toronto</town><town>Ottawa</town></towns>');
```

The `BY REF` clauses have no effect in PostgreSQL, but are allowed for SQL conformance and compatibility with other implementations. Per SQL standard, the first `BY REF` is required, the second is optional. Also note that the SQL standard specifies the `xmlexists` construct to take an XQuery expression as first argument, but PostgreSQL currently only supports XPath, which is a subset of XQuery.

### 9.14.2.4. `xml_is_well_formed`

These functions check whether a `text` string is well-formed XML, returning a Boolean result. `xml_is_well_formed_document` checks for a well-formed document, while `xml_is_well_formed_content` checks for well-formed content. `xml_is_well_formed` does the former if the `xmloption` configuration parameter is set to `DOCUMENT`, or the latter if it is set to `CONTENT`. This means that `xml_is_well_formed` is useful for seeing whether a simple cast to type `xml` will succeed, whereas the other two functions are useful for seeing whether the corresponding variants of `XMLPARSE` will succeed.

Examples:

```sql
SET xmloption TO DOCUMENT;
SELECT xml_is_well_formed('<>');
```

```
8
```

```sql
SELECT xml_is_well_formed('<abc/>');
```

```
8
```

```sql
SET xmloption TO CONTENT;
SELECT xml_is_well_formed('abc');
```

```
8
```
The last example shows that the checks include whether namespaces are correctly matched.

### 9.14.3. Processing XML

To process values of data type `xml`, PostgreSQL offers the functions `xpath` and `xpath_exists`, which evaluate XPath 1.0 expressions, and the `XMLTABLE` table function.

#### 9.14.3.1. `xpath`

`xpath(xpath, xml [, nsarray])`

The function `xpath` evaluates the XPath expression `xpath(a text value)` against the XML value `xml`. It returns an array of XML values corresponding to the node set produced by the XPath expression. If the XPath expression returns a scalar value rather than a node set, a single-element array is returned.

The second argument must be a well formed XML document. In particular, it must have a single root node element.

The optional third argument of the function is an array of namespace mappings. This array should be a two-dimensional `text` array with the length of the second axis being equal to 2 (i.e., it should be an array of arrays, each of which consists of exactly 2 elements). The first element of each array entry is the namespace name (alias), the second the namespace URI. It is not required that aliases provided in this array be the same as those being used in the XML document itself (in other words, both in the XML document and in the `xpath` function context, aliases are `local`).

Example:

```sql
```

```plaintext
xpath
-------
{test}
(1 row)
```

To deal with default (anonymous) namespaces, do something like this:
9.14.3.2. xpath_exists

 xpath_exists (xpath, xml [, nsarray])

The function xpath_exists is a specialized form of the xpath function. Instead of returning the individual XML values that satisfy the XPath, this function returns a Boolean indicating whether the query was satisfied or not. This function is equivalent to the standard XMLEXISTS predicate, except that it also offers support for a namespace mapping argument.

Example:

ARRAY[ARRAY['my', 'http://example.com']] );

xpath_exists
-----------------
 t
(1 row)

9.14.3.3. xmltable

xmltable( [XMLNAMESPACES(namespace uri AS namespace name[, ...]), ]
row_expression PASSING [BY REF] document_expression [BY REF]

COLUMNS name { type [PATH column_expression]
[DEFAULT default_expression] [NOT NULL | NULL]
[ | FOR ORDINALITY ]
[ , ... ]
)

The xmltable function produces a table based on the given XML value, an XPath filter to extract rows, and an optional set of column definitions.

The optional XMLNAMESPACES clause is a comma-separated list of namespaces. It specifies the XML namespaces used in the document and their aliases. A default namespace specification is not currently supported.

The required row_expression argument is an XPath expression that is evaluated against the supplied XML document to obtain an ordered sequence of XML nodes. This sequence is what xmltable transforms into output rows.

document_expression provides the XML document to operate on. The BY REF clauses have no effect in PostgreSQL, but are allowed for SQL conformance and compatibility with other implementations. The argument must be a well-formed XML document; fragments/forests are not accepted.
The mandatory **COLUMNS** clause specifies the list of columns in the output table. If the **COLUMNS** clause is omitted, the rows in the result set contain a single column of type **xml** containing the data matched by **row_expression**. If **COLUMNS** is specified, each entry describes a single column. See the syntax summary above for the format. The column name and type are required; the path, default and nullability clauses are optional.

A column marked **FOR ORDINALITY** will be populated with row numbers matching the order in which the output rows appeared in the original input XML document. At most one column may be marked **FOR ORDINALITY**.

The **column_expression** for a column is an XPath expression that is evaluated for each row, relative to the result of the **row_expression**, to find the value of the column. If no **column_expression** is given, then the column name is used as an implicit path.

If a column's XPath expression returns multiple elements, an error is raised. If the expression matches an empty tag, the result is an empty string (not **NULL**). Any **xsi:nil** attributes are ignored.

The text body of the XML matched by the **column_expression** is used as the column value. Multiple **text()** nodes within an element are concatenated in order. Any child elements, processing instructions, and comments are ignored, but the text contents of child elements are concatenated to the result. Note that the whitespace-only **text()** node between two non-text elements is preserved, and that leading whitespace on a **text()** node is not flattened.

If the path expression does not match for a given row but **default_expression** is specified, the value resulting from evaluating that expression is used. If no **DEFAULT** clause is given for the column, the field will be set to **NULL**. It is possible for a **default_expression** to reference the value of output columns that appear prior to it in the column list, so the default of one column may be based on the value of another column.

Columns may be marked **NOT NULL**. If the **column_expression** for a **NOT NULL** column does not match anything and there is no **DEFAULT** or the **default_expression** also evaluates to null, an error is reported.

Unlike regular PostgreSQL functions, **column_expression** and **default_expression** are not evaluated to a simple value before calling the function. **column_expression** is normally evaluated exactly once per input row, and **default_expression** is evaluated each time a default is needed for a field. If the expression qualifies as stable or immutable the repeat evaluation may be skipped. Effectively **xmltable** behaves more like a subquery than a function call. This means that you can usefully use volatile functions like **nextval** in **default_expression** and **column_expression** may depend on other parts of the XML document.

Examples:

```sql
CREATE TABLE xmldata AS SELECT
  xml $$
  <ROWS>
    <ROW id="1">
      <COUNTRY_ID>AU</COUNTRY_ID>
      <COUNTRY_NAME>Australia</COUNTRY_NAME>
    </ROW>
    <ROW id="5">
      <COUNTRY_ID>JP</COUNTRY_ID>
      <COUNTRY_NAME>Japan</COUNTRY_NAME>
      <PREMIER_NAME>Shinzo Abe</PREMIER_NAME>
      <SIZE unit="sq_mi">145935</SIZE>
    </ROW>
    <ROW id="6">
      <COUNTRY_ID>SG</COUNTRY_ID>
  </ROWS>
$$;
```
The following example shows concatenation of multiple text() nodes, usage of the column name as XPath filter, and the treatment of whitespace, XML comments and processing instructions:

```
CREATE TABLE xmlelements AS SELECT
xml $$
<root>
  <element> Hello<!-- xyxxz -->2a2<?aaaaa?> <!--x--> bbb<x>xxx</x>CC </element>
</root>
$$ AS data;

SELECT xmltable.*
FROM xmlelements,
XMLTABLE('/root' PASSING data COLUMNS element
  element
) ;
```

```
| element |
Hello2a2 bbbCC
```

The following example illustrates how the XMLNAMESPACES clause can be used to specify a list of namespaces used in the XML document as well as in the XPath expressions:

```
WITH xmldata(data) AS (VALUES ('
<example xmlns="http://example.com/myns" xmlns:B="http://example.com/b"

SELECT xmltable.*
FROM xmldata,
XMLTABLE('/example' PASSING data COLUMNS element
  element
) ;
```

```
| element |
Hello2a2 bbbCC
```
Functions and Operators

```sql
SELECT xmltable.*
FROM XMLTABLE(XMLNAMESPACES('http://example.com/myns' AS x, 'http://example.com/b' AS "B"),
'/x:example/x:item'
PASSING (SELECT data FROM xmldata)
COLUMNS foo int PATH '@foo',
bar int PATH '@B:bar');
```

<table>
<thead>
<tr>
<th>foo</th>
<th>bar</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2</td>
</tr>
<tr>
<td>3</td>
<td>4</td>
</tr>
<tr>
<td>4</td>
<td>5</td>
</tr>
</tbody>
</table>

(3 rows)

9.14.4. Mapping Tables to XML

The following functions map the contents of relational tables to XML values. They can be thought of as XML export functionality:

```sql
table_to_xml(tbl regclass, nulls boolean, tableforest boolean, targetns text)
query_to_xml(query text, nulls boolean, tableforest boolean, targetns text)
cursor_to_xml(cursor refcursor, count int, nulls boolean, tableforest boolean, targetns text)
```

The return type of each function is `xml`.

`table_to_xml` maps the content of the named table, passed as parameter `tbl`. The `regclass` type accepts strings identifying tables using the usual notation, including optional schema qualifications and double quotes. `query_to_xml` executes the query whose text is passed as parameter `query` and maps the result set. `cursor_to_xml` fetches the indicated number of rows from the cursor specified by the parameter `cursor`. This variant is recommended if large tables have to be mapped, because the result value is built up in memory by each function.

If `tableforest` is false, then the resulting XML document looks like this:

```xml
<tablename>
  <row>
    <columnname1>data</columnname1>
    <columnname2>data</columnname2>
  </row>
  ...
</tablename>
```

If `tableforest` is true, the result is an XML content fragment that looks like this:


Functions and Operators

If no table name is available, that is, when mapping a query or a cursor, the string table is used in the first format, row in the second format.

The choice between these formats is up to the user. The first format is a proper XML document, which will be important in many applications. The second format tends to be more useful in the cursor_to_xml function if the result values are to be reassembled into one document later on. The functions for producing XML content discussed above, in particular xmlelement, can be used to alter the results to taste.

The data values are mapped in the same way as described for the function xmlelement above.

The parameter nulls determines whether null values should be included in the output. If true, null values in columns are represented as:

\[
<\text{columnname} \text{xsi:nil="true"} />
\]

where xsi is the XML namespace prefix for XML Schema Instance. An appropriate namespace declaration will be added to the result value. If false, columns containing null values are simply omitted from the output.

The parameter targetns specifies the desired XML namespace of the result. If no particular namespace is wanted, an empty string should be passed.

The following functions return XML Schema documents describing the mappings performed by the corresponding functions above:

\[
\text{table_to_xmlschema} (\text{tbl regclass, nulls boolean, tableforest boolean, targetns text})
\]
\[
\text{query_to_xmlschema} (\text{query text, nulls boolean, tableforest boolean, targetns text})
\]
\[
\text{cursor_to_xmlschema} (\text{cursor refcursor, nulls boolean, tableforest boolean, targetns text})
\]

It is essential that the same parameters are passed in order to obtain matching XML data mappings and XML Schema documents.

The following functions produce XML data mappings and the corresponding XML Schema in one document (or forest), linked together. They can be useful where self-contained and self-describing results are wanted:

\[
\text{table_to_xml_and_xmlschema} (\text{tbl regclass, nulls boolean, tableforest boolean, targetns text})
\]
\[
\text{query_to_xml_and_xmlschema} (\text{query text, nulls boolean, tableforest boolean, targetns text})
\]
In addition, the following functions are available to produce analogous mappings of entire schemas or the entire current database:

\begin{verbatim}
schema_to_xml(schema name, nulls boolean, tableforest boolean, targetns text)
schema_to_xmlschema(schema name, nulls boolean, tableforest boolean, targetns text)
schema_to_xml_and_xmlschema(schema name, nulls boolean, tableforest boolean, targetns text)
database_to_xml(nulls boolean, tableforest boolean, targetns text)
database_to_xmlschema(nulls boolean, tableforest boolean, targetns text)
database_to_xml_and_xmlschema(nulls boolean, tableforest boolean, targetns text)
\end{verbatim}

Note that these potentially produce a lot of data, which needs to be built up in memory. When requesting content mappings of large schemas or databases, it might be worthwhile to consider mapping the tables separately instead, possibly even through a cursor.

The result of a schema content mapping looks like this:

\begin{verbatim}
<schemaname>
table1-mapping
table2-mapping...
</schemaname>
\end{verbatim}

where the format of a table mapping depends on the `tableforest` parameter as explained above.

The result of a database content mapping looks like this:

\begin{verbatim}
<dbname>
<schema1name>...
</schema1name>
<schema2name>...
</schema2name>...
</dbname>
\end{verbatim}

where the schema mapping is as above.

As an example of using the output produced by these functions, Figure 9.1 shows an XSLT stylesheet that converts the output of `table_to_xml_and_xmlschema` to an HTML document containing a tabular rendition of the table data. In a similar manner, the results from these functions can be converted into other XML-based formats.
Figure 9.1. XSLT Stylesheet for Converting SQL/XML Output to HTML

```xml
<?xml version="1.0"?
<xsl:stylesheet version="1.0"
 xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
 xmlns:xsd="http://www.w3.org/2001/XMLSchema"
 xmlns="http://www.w3.org/1999/xhtml">
  <xsl:output method="xml"
    doctype-system="http://www.w3.org/TR/xhtml1/DTD/xhtml1-strict.dtd"
    doctype-public="-//W3C/DTD XHTML 1.0 Strict//EN"
    indent="yes"/>

  <xsl:template match="/*">
    <xsl:variable name="schema" select="//xsd:schema"/>
    <xsl:variable name="tabletypename" select="$schema/xsd:element[@name='name(current())']/@type"/>
    <xsl:variable name="rowtypename" select="$schema/xsd:complexType[@name='$tabletypename']/xsd:sequence/xsd:element[@name='row']/@type"/>

    <html>
      <head>
        <title><xsl:value-of select="name(current())"/></title>
      </head>
      <body>
        <table>
          <tr>
            <xsl:for-each select="$schema/xsd:complexType[@name='$rowtypename']/xsd:sequence/xsd:element[@name='name']/@name">
              <th><xsl:value-of select="."/></th>
            </xsl:for-each>
          </tr>
          <xsl:for-each select="row">
            <tr>
              <xsl:for-each select="*">
                <td><xsl:value-of select="."/></td>
              </xsl:for-each>
            </tr>
          </xsl:for-each>
        </table>
      </body>
    </html>

  </xsl:template>
</xsl:stylesheet>
```

9.15. JSON Functions and Operators

Table 9.43 shows the operators that are available for use with the two JSON data types (see Section 8.14).
Table 9.43. `json` and `jsonb` Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Right Type</th>
<th>Operand</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>-&gt;</code></td>
<td>int</td>
<td>Get JSON array <code>'{&quot;a&quot;:&quot;foo&quot;,&quot;c&quot;:&quot;baz&quot;}'</code> element (indexed by zero, negative integers count from the end)</td>
<td><code>son-&gt;'a'</code></td>
<td><code>{&quot;c&quot;:&quot;baz&quot;}</code></td>
<td></td>
</tr>
<tr>
<td><code>-&gt;</code></td>
<td>text</td>
<td>Get JSON object <code>'{&quot;a&quot;:&quot;foo&quot;,&quot;b&quot;:&quot;baz&quot;}'</code> field by key</td>
<td><code>son-&gt;'a'</code></td>
<td><code>{&quot;b&quot;:&quot;baz&quot;}</code></td>
<td></td>
</tr>
<tr>
<td><code>-&gt;&gt;</code></td>
<td>int</td>
<td>Get JSON array element as text</td>
<td><code>son-&gt;2</code></td>
<td><code>{&quot;c&quot;:&quot;baz&quot;}</code></td>
<td></td>
</tr>
<tr>
<td><code>-&gt;&gt;</code></td>
<td>text</td>
<td>Get JSON object field as text</td>
<td><code>son-&gt;'a'</code></td>
<td><code>{&quot;b&quot;:&quot;baz&quot;}</code></td>
<td></td>
</tr>
<tr>
<td><code>#&gt;</code></td>
<td>text[]</td>
<td>Get JSON object at specified path</td>
<td><code>son#&gt;'{a,b}'</code></td>
<td><code>{&quot;c&quot;: &quot;baz&quot;}</code></td>
<td></td>
</tr>
<tr>
<td><code>#&gt;&gt;</code></td>
<td>text[]</td>
<td>Get JSON object at specified path as text</td>
<td><code>son#&gt;&gt;'{a,2}'</code></td>
<td><code>3</code></td>
<td></td>
</tr>
</tbody>
</table>

Note

There are parallel variants of these operators for both the `json` and `jsonb` types. The field/element/path extraction operators return the same type as their left-hand input (either `json` or `jsonb`), except for those specified as returning `text`, which coerce the value to `text`. The field/element/path extraction operators return NULL, rather than failing, if the JSON input does not have the right structure to match the request; for example if no such element exists. The field/element/path extraction operators that accept integer JSON array subscripts all support negative subscripts from the end of arrays.

The standard comparison operators shown in Table 9.1 are available for `jsonb`, but not for `json`. They follow the ordering rules for B-tree operations outlined at Section 8.14.4.

Some further operators also exist only for `jsonb`, as shown in Table 9.44. Many of these operators can be indexed by `jsonb` operator classes. For a full description of `jsonb` containment and existence semantics, see Section 8.14.3. Section 8.14.4 describes how these operators can be used to effectively index `jsonb` data.

Table 9.44. Additional `jsonb` Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Right Operand Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>@&gt;</code></td>
<td><code>jsonb</code></td>
<td>Does the left JSON value contain the right <code>&quot;b&quot;:2'</code> at the top level?</td>
<td><code>sonb</code></td>
</tr>
<tr>
<td><code>&lt;@</code></td>
<td><code>jsonb</code></td>
<td>Are the left JSON path/value entries contained</td>
<td><code>sonb</code></td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Right Operand Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>?</td>
<td>text</td>
<td>Does the string exist as a top-level key within the JSON value?</td>
<td>&quot;b&quot;:2::jsonb ? 'b'</td>
</tr>
<tr>
<td>?</td>
<td></td>
<td>text[]</td>
<td>Do any of these strings exist as top-level keys?</td>
</tr>
<tr>
<td>?&amp;</td>
<td>text[]</td>
<td>Do all of these strings exist as top-level keys?</td>
<td>sonb ?&amp; array['a', 'b']</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>jsonb</td>
</tr>
<tr>
<td>-</td>
<td>text</td>
<td>Delete key/value pair or string element from left operand. Key/value pairs are matched on their key value.</td>
<td>&quot;a&quot;: &quot;b&quot;</td>
</tr>
<tr>
<td>-</td>
<td>text[]</td>
<td>Delete multiple key/ value pairs or string elements from left operand. Key/value pairs are matched on their key value.</td>
<td>&quot;a&quot;, &quot;b&quot;, &quot;c&quot;:&quot;d&quot;</td>
</tr>
<tr>
<td>-</td>
<td>integer</td>
<td>Delete the array element with specified index (Negative integers count from the end). Throws an error if top level container is not an array.</td>
<td>'a&quot;, &quot;b&quot;:1</td>
</tr>
<tr>
<td>#-</td>
<td>text[]</td>
<td>Delete the field or element with specified path (for JSON arrays, negative integers count from the end).</td>
<td>&quot;a&quot;:1</td>
</tr>
</tbody>
</table>

| Note |

The operator concatenates the elements at the top level of each of its operands. It does not operate recursively. For example, if both operands are objects with a common key field name, the value of the field in the result will just be the value from the right hand operand.

Table 9.45 shows the functions that are available for creating json and jsonb values. (There are no equivalent functions for jsonb, of the row_to_json and array_to_json functions. However, the to_jsonb function supplies much the same functionality as these functions would.)
### Table 9.45. JSON Creation Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
</table>
| to_json(anyelement)     | Returns the value as `json` or `jsonb`. Arrays and composites are converted (recursively) to arrays and objects; otherwise, if there is a cast from the type to `json`, the cast function will be used to perform the conversion; otherwise, a scalar value is produced. For any scalar type other than a number, a Boolean, or a null value, the text representation will be used, in such a fashion that it is a valid `json` or `jsonb` value. | to_json('Fred said "Hi."::text)          | "Fred said "Hi."
                       |                                        |                                          | ""                  |
| to_jsonb(anyelement)    |                                                                                                |                                          |                    |
| array_to_json(anyarray [, pretty_bool]) | Returns the array as a JSON array. A PostgreSQL multidimensional array becomes a JSON array of arrays. Line feeds will be added between dimension-1 elements if `pretty_bool` is true. | array_to_json('{{1,5},{99,100}}'::int[]) | [[1,5],[99,100]] |
| row_to_json(record [, pretty_bool]) | Returns the row as a JSON object. Line feeds will be added between level-1 elements if `pretty_bool` is true. | row_to_json(row(1,'foo'))             | "f1": 1,"f2": "foo" |
| json_build_array(VARIADIC "any") | Builds a possibly-heterogeneously-typed JSON array out of a variadic argument list.                                                         | json_build_array(1,2,'3',4,5)          | [1, 2, "3", 4, 5] |
| jsonb_build_array(VARIADIC "any") |                                                                                                |                                          |                    |
| json_build_object(VARIADIC "any") | Builds a JSON object out of a variadic argument list. By convention, the argument list consists of alternating keys and values. | json_build_object("foo": 1, "bar": 2) |                     |
| jsonb_build_object(VARIADIC "any") |                                                                                                |                                          |                    |
| json_object(text[])     | Builds a JSON object out of a text array. The array must have either exactly one dimension with an even number of keys.                  | json_object(‘a, {"a": "1", "b": 1, b, "def", c, "def", "c": 3.5}’) |                     |
| jsonb_object(text[])    |                                                                                                |                                          |                    |
Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>members, in which case</td>
<td>they are taken as alternating key/value pairs, or two dimensions such that each inner array has exactly two elements, which are taken as a key/value pair.</td>
<td><code>json_object('{{a, 1}, {b, &quot;def&quot;}, {c, 3.5}}')</code>)</td>
<td></td>
</tr>
<tr>
<td>json_object(keys text[],</td>
<td>This form of json_object takes keys and values pairwise from two separate arrays. In all other respects it is identical to the one-argument form.</td>
<td><code>json_object('{{a, &quot;1&quot;}, {b: &quot;2&quot;, c: &quot;3.5&quot;}}')</code>)</td>
<td></td>
</tr>
<tr>
<td>values text[])</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>jsonb_object(keys text[],</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>values text[])</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Note

array_to_json and row_to_json have the same behavior as to_json except for offering a pretty-printing option. The behavior described for to_json likewise applies to each individual value converted by the other JSON creation functions.

Note

The hstore extension has a cast from hstore to json, so that hstore values converted via the JSON creation functions will be represented as JSON objects, not as primitive string values.

Table 9.46 shows the functions that are available for processing json and jsonb values.

Table 9.46. JSON Processing Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>json_array_length(json)</td>
<td>int</td>
<td>Returns the number of elements in the outermost JSON array.</td>
<td><code>json_array_length('[[1,2,3],[&quot;f1&quot;:1,&quot;f2&quot;:2]]')</code>)</td>
<td>5</td>
</tr>
<tr>
<td>jsonb_array_length(jsonb)</td>
<td>setof</td>
<td>Expands the outermost JSON object into a set of key/value pairs.</td>
<td><code>jsonb_array_length('&quot;a&quot;:&quot;foo&quot;, &quot;b&quot;:&quot;bar&quot;')</code>)</td>
<td>+-------</td>
</tr>
<tr>
<td>json_each(json text,</td>
<td>setof</td>
<td>Expands the outermost JSON object from key/value pairs.</td>
<td><code>json_each(text, value)</code>)</td>
<td>a</td>
</tr>
<tr>
<td>json each(json text,</td>
<td>setof</td>
<td>Expands the outermost JSON object from key/value pairs.</td>
<td><code>json each('&quot;a&quot;:&quot;foo&quot;, &quot;b&quot;:&quot;bar&quot;')</code>)</td>
<td>a</td>
</tr>
<tr>
<td>jsonb_each(json text,</td>
<td>setof</td>
<td>Expands the outermost JSON object from key/value pairs.</td>
<td><code>jsonb each('&quot;a&quot;:&quot;foo&quot;, &quot;b&quot;:&quot;bar&quot;')</code>)</td>
<td>a</td>
</tr>
<tr>
<td>json_each_text setof</td>
<td>setof</td>
<td>Expands the outermost JSON object from key/value pairs.</td>
<td><code>json_each_text setof('&quot;a&quot;:&quot;foo&quot;, &quot;b&quot;:&quot;bar&quot;')</code>)</td>
<td>a</td>
</tr>
<tr>
<td>jsonb_each_text setof</td>
<td>setof</td>
<td>Expands the outermost JSON object from key/value pairs.</td>
<td><code>jsonb_each_text setof('&quot;a&quot;:&quot;foo&quot;, &quot;b&quot;:&quot;bar&quot;')</code>)</td>
<td>a</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Example Result</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
<td>---------</td>
<td>---------------</td>
</tr>
<tr>
<td>jsonb_each_text(jsonb)</td>
<td></td>
<td>turned values will be of type text.</td>
<td></td>
<td>b</td>
</tr>
<tr>
<td>json_extract_path(from_json json, VARIADIC jsonb, path_elems text[])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>jsonb_extract_path(from_json jsonb, VARIADIC path_elems text[])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_extract_path_text(from_json json, VARIADIC path_elems text[])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>jsonb_extract_path_text(from_json jsonb, VARIADIC path_elems text[])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_object_keys(json)</td>
<td>setof text</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>jsonb_object_keys(jsonb)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_populate_record(base anyelement, from_json json)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>jsonb_populate_record(base anyelement, from_json jsonb)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_populate_record_set(base anyelement, setof anyelement from_json json)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_populate_record_set(base anyelement, setof anyelement from_json jsonb)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Example Result</td>
</tr>
<tr>
<td>---------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>-------------------------------------------------------------------------</td>
<td>-------------------------------------</td>
</tr>
<tr>
<td>from_json json</td>
<td></td>
<td>columns match the record type defined by base (see note below).</td>
<td><code>{&quot;a&quot;:1,&quot;b&quot;:2}, {&quot;a&quot;:3,&quot;b&quot;:4}</code></td>
<td></td>
</tr>
<tr>
<td>jsonb_populate_recordset</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_arrray_elements(json)</td>
<td>setof json</td>
<td>Expands a JSON array to a set of JSON values.</td>
<td>select * from json_arrray_elements([1,true],[2,false])</td>
<td>true [2,false]</td>
</tr>
<tr>
<td></td>
<td>setof jsonb</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_arrray_elements_text(json)</td>
<td>setof text</td>
<td>Expands a JSON array to a set of text values.</td>
<td>select * from json_arrray_elements_text('&quot;foo&quot;,&quot;bar&quot;')</td>
<td>foo bar</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_typeof(json)</td>
<td>text</td>
<td>Returns the type of the outermost JSON value as a text string.</td>
<td>json_typeof('-123.4')</td>
<td>number</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_to_record(json)</td>
<td></td>
<td>Builds an arbitrary record from a JSON object.</td>
<td>json_to_record('{{&quot;a&quot;:1,&quot;b&quot;:2,&quot;c&quot;:&quot;foo&quot;,&quot;d&quot;:&quot;bar&quot;,&quot;r&quot;:&quot;baz&quot;}}')</td>
<td>x(a int, b text, c string, d text,</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>r myrowtype)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_to_record_of_record</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>set(json)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_arrray_elements_set(json)</td>
<td>setof record</td>
<td>Builds an arbitrary record set from a JSON array of objects.</td>
<td>json_to_record_of_record_set('{{&quot;a&quot;:1,&quot;b&quot;:&quot;foo&quot;,&quot;c&quot;:&quot;bar&quot;}}')</td>
<td>x(a int, b text, c string)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>json_arrray_elements_text_set(json)</td>
<td>setof</td>
<td>Builds an arbitrary record set from a JSON array of objects.</td>
<td>json_to_record_of_record_set('{{&quot;a&quot;:1,&quot;b&quot;:&quot;foo&quot;,&quot;c&quot;:&quot;bar&quot;}}')</td>
<td>x(a int, b text, c string)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Functions and Operators

Function

Return Type

Description
Example
plicitly define the
structure of the
record with an AS
clause.

Example Result

Returns from_json with all object fields that have
null values omitted. Other null values are untouched.

json_strip_null[{"f1":1},2,nuls('[{"f1":1,"f2":null,3]
l},2,null,3]')

jsonjsonb
b_set(target
jsonb,
path
text[],
new_value
jsonb[, create_missing
boolean])

Returns target
with
the
section designated by
path replaced by
new_value, or
with new_value added if create_missing is
true ( default is
true) and the
item designated by
path does not
exist. As with
the path orientated operators, negative integers that
appear in path
count from the end
of JSON arrays.

json[{"f1":
b_set('[{"f1":1,"f2":nul[2,3,4],"f2":null},2,null},2,null,3]
l,3]',
[{"f1":
1,
'{0,f1}','[2,3,4]',
"f2":
null,
false)
"f3": [2, 3,
json4]}, 2]
b_set('[{"f1":1,"f2":null},2]',
'{0,f3}','[2,3,4]')

jsonb_injsonb
sert(target
jsonb,
path
text[],
new_value
jsonb,
[insert_after
boolean])

Returns target
with new_value
inserted. If target section designated by path
is in a JSONB
array, new_value will be inserted before target or after if insert_after is
true (default is
false). If target section designated by path
is in JSONB object, new_value
will be inserted only if target does
not exist. As with
the path orientated operators, negative integers that
appear in path

jsonb_insert('{"a":
[0,1,2]}',
'{a,
1}',
'"new_value"')

json_strip_nulljson
s(from_json
jsonb
json)
jsonb_strip_nulls(from_json
jsonb)

291

jsonb_insert('{"a":
[0,1,2]}',
'{a,
1}',
'"new_value"', true)

{"a":
[0,
"new_value",
1, 2]}
{"a":
[0,
1, "new_value", 2]}


### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>jsonb_pretty_text</td>
<td>text</td>
<td>Returns <code>from_json</code> as indented JSON text.</td>
<td><code>jsonb_pretty_text('&quot;\{&quot;\&quot;f1\&quot;:1,\&quot;f2\&quot;:null,\&quot;f3\&quot;:3\&quot;\}', \{&quot;f1&quot;: 1, \&quot;f2&quot;: null, 2, null, 3\})</code></td>
<td><code>\{&quot;f1&quot;: 1, \&quot;f2&quot;: null, 2, null, 3\}</code></td>
</tr>
</tbody>
</table>

**Note**

Many of these functions and operators will convert Unicode escapes in JSON strings to the appropriate single character. This is a non-issue if the input is type `jsonb`, because the conversion was already done; but for `json` input, this may result in throwing an error, as noted in Section 8.14.

**Note**

While the examples for the functions `json_populate_record`, `json_populate_recordset`, `json_to_record` and `json_to_recordset` use constants, the typical use would be to reference a table in the `FROM` clause and use one of its `json` or `jsonb` columns as an argument to the function. Extracted key values can then be referenced in other parts of the query, like `WHERE` clauses and target lists. Extracting multiple values in this way can improve performance over extracting them separately with per-key operators.

JSON keys are matched to identical column names in the target row type. JSON type coercion for these functions is “best effort” and may not result in desired values for some types. JSON fields that do not appear in the target row type will be omitted from the output, and target columns that do not match any JSON field will simply be NULL.

**Note**

All the items of the `path` parameter of `jsonb_set` as well as `jsonb_insert` except the last item must be present in the target. If `create_missing` is false, all items of the `path` parameter of `jsonb_set` must be present. If these conditions are not met the target is returned unchanged.

If the last path item is an object key, it will be created if it is absent and given the new value. If the last path item is an array index, if it is positive the item to set is found by counting from the left, and if negative by counting from the right -1 designates the rightmost element, and so on. If the item is out of the range `array_length .. array_length -1`, and `create_missing` is true, the new value is added at the beginning of the array if the item is negative, and at the end of the array if it is positive.
Note

The \texttt{json_typeof} function's \texttt{null} return value should not be confused with a SQL NULL. While calling \texttt{json_typeof('null':::json)} will return \texttt{null}, calling \texttt{json_typeof(NULL::json)} will return a SQL NULL.

Note

If the argument to \texttt{json_strip_nulls} contains duplicate field names in any object, the result could be semantically somewhat different, depending on the order in which they occur. This is not an issue for \texttt{jsonb_strip_nulls} since \texttt{jsonb} values never have duplicate object field names.

See also Section~9.20 for the aggregate function \texttt{json_agg} which aggregates record values as JSON, and the aggregate function \texttt{json_object_agg} which aggregates pairs of values into a JSON object, and their \texttt{jsonb} equivalents, \texttt{jsonb_agg} and \texttt{jsonb_object_agg}.

9.16. Sequence Manipulation Functions

This section describes functions for operating on \emph{sequence objects}, also called sequence generators or just sequences. Sequence objects are special single-row tables created with \texttt{CREATE SEQUENCE}. Sequence objects are commonly used to generate unique identifiers for rows of a table. The sequence functions, listed in Table~9.47, provide simple, multiuser-safe methods for obtaining successive sequence values from sequence objects.

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{currval(regclass)}</td>
<td>\texttt{bigint}</td>
<td>Return value most recently obtained with \texttt{nextval} for specified sequence</td>
</tr>
<tr>
<td>\texttt{lastval()}</td>
<td>\texttt{bigint}</td>
<td>Return value most recently obtained with \texttt{nextval} for any sequence</td>
</tr>
<tr>
<td>\texttt{nextval(regclass)}</td>
<td>\texttt{bigint}</td>
<td>Advance sequence and return new value</td>
</tr>
<tr>
<td>\texttt{setval(regclass, bigint)}</td>
<td>\texttt{bigint}</td>
<td>Set sequence's current value</td>
</tr>
<tr>
<td>\texttt{setval(regclass, bigint, boolean)}</td>
<td>\texttt{bigint}</td>
<td>Set sequence's current value and is_called flag</td>
</tr>
</tbody>
</table>

The sequence to be operated on by a sequence function is specified by a \texttt{regclass} argument, which is simply the OID of the sequence in the \texttt{pg\_class} system catalog. You do not have to look up the OID by hand, however, since the \texttt{regclass} data type's input converter will do the work for you. Just write the sequence name enclosed in single quotes so that it looks like a literal constant. For compatibility with the handling of ordinary SQL names, the string will be converted to lower case unless it contains double quotes around the sequence name. Thus:

\begin{verbatim}
nextval('foo')  \hspace{1cm} operates on sequence foo
nextval('FOO')  \hspace{1cm} operates on sequence foo
nextval("Foo") \hspace{1cm} operates on sequence Foo
\end{verbatim}
The sequence name can be schema-qualified if necessary:

nextval('myschema.foo') operates on myschema.foo
nextval('"myschema".foo') same as above
nextval('foo') searches search path for foo

See Section 8.19 for more information about regclass.

Note

Before PostgreSQL 8.1, the arguments of the sequence functions were of type text, not regclass, and the above-described conversion from a text string to an OID value would happen at run time during each call. For backward compatibility, this facility still exists, but internally it is now handled as an implicit coercion from text to regclass before the function is invoked.

When you write the argument of a sequence function as an unadorned literal string, it becomes a constant of type regclass. Since this is really just an OID, it will track the originally identified sequence despite later renaming, schema reassignment, etc. This "early binding" behavior is usually desirable for sequence references in column defaults and views. But sometimes you might want "late binding" where the sequence reference is resolved at run time. To get late-binding behavior, force the constant to be stored as a text constant instead of regclass:

nextval('foo':::text) foo is looked up at runtime

Note that late binding was the only behavior supported in PostgreSQL releases before 8.1, so you might need to do this to preserve the semantics of old applications.

Of course, the argument of a sequence function can be an expression as well as a constant. If it is a text expression then the implicit coercion will result in a run-time lookup.

The available sequence functions are:

nextval

Advance the sequence object to its next value and return that value. This is done atomically: even if multiple sessions execute nextval concurrently, each will safely receive a distinct sequence value.

If a sequence object has been created with default parameters, successive nextval calls will return successive values beginning with 1. Other behaviors can be obtained by using special parameters in the CREATE SEQUENCE command; see its command reference page for more information.

Important

To avoid blocking concurrent transactions that obtain numbers from the same sequence, a nextval operation is never rolled back; that is, once a value has been fetched it is considered used and will not be returned again. This is true even if the surrounding transaction later aborts, or if the calling query ends up not using the value. For example an INSERT with an ON CONFLICT clause will compute the to-be-inserted tuple, including doing any required nextval calls, before detecting any conflict that would cause it to follow the ON CONFLICT rule instead.
Such cases will leave unused “holes” in the sequence of assigned values. Thus, PostgreSQL sequence objects cannot be used to obtain “gapless” sequences.

This function requires USAGE or UPDATE privilege on the sequence.

**currval**

Return the value most recently obtained by `nextval` for this sequence in the current session. (An error is reported if `nextval` has never been called for this sequence in this session.) Because this is returning a session-local value, it gives a predictable answer whether or not other sessions have executed `nextval` since the current session did.

This function requires USAGE or SELECT privilege on the sequence.

**lastval**

Return the value most recently returned by `nextval` in the current session. This function is identical to `currval`, except that instead of taking the sequence name as an argument it refers to whichever sequence `nextval` was most recently applied to in the current session. It is an error to call `lastval` if `nextval` has not yet been called in the current session.

This function requires USAGE or SELECT privilege on the last used sequence.

**setval**

Reset the sequence object's counter value. The two-parameter form sets the sequence's `last_value` field to the specified value and sets its `is_called` field to `true`, meaning that the next `nextval` will advance the sequence before returning a value. The value reported by `currval` is also set to the specified value. In the three-parameter form, `is_called` can be set to either `true` or `false`. `true` has the same effect as the two-parameter form. If it is set to `false`, the next `nextval` will return exactly the specified value, and sequence advancement commences with the following `nextval`. Furthermore, the value reported by `currval` is not changed in this case. For example,

```
SELECT setval('foo', 42);         Next nextval will return 43
SELECT setval('foo', 42, true);  Same as above
SELECT setval('foo', 42, false); Next nextval will return 42
```

The result returned by `setval` is just the value of its second argument.

**Important**

Because sequences are non-transactional, changes made by `setval` are not undone if the transaction rolls back.

This function requires UPDATE privilege on the sequence.

### 9.17. Conditional Expressions

This section describes the SQL-compliant conditional expressions available in PostgreSQL.

**Tip**

If your needs go beyond the capabilities of these conditional expressions, you might want to consider writing a server-side function in a more expressive programming language.
9.17.1. CASE

The SQL **CASE** expression is a generic conditional expression, similar to if/else statements in other programming languages:

```
CASE WHEN condition THEN result
    [WHEN ...]
    [ELSE result]
END
```

**CASE** clauses can be used wherever an expression is valid. Each **condition** is an expression that returns a **boolean** result. If the condition's result is true, the value of the **CASE** expression is the **result** that follows the condition, and the remainder of the **CASE** expression is not processed. If the condition's result is not true, any subsequent **WHEN** clauses are examined in the same manner. If no **WHEN condition** yields true, the value of the **CASE** expression is the **result** of the **ELSE** clause. If the **ELSE** clause is omitted and no condition is true, the result is null.

An example:

```
SELECT * FROM test;

<table>
<thead>
<tr>
<th>a</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
</tr>
<tr>
<td>2</td>
</tr>
<tr>
<td>3</td>
</tr>
</tbody>
</table>
```

```
SELECT a,
    CASE WHEN a=1 THEN 'one'
        WHEN a=2 THEN 'two'
        ELSE 'other'
    END
FROM test;

<table>
<thead>
<tr>
<th>a</th>
<th>case</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
</tr>
<tr>
<td>3</td>
<td>other</td>
</tr>
</tbody>
</table>
```

The data types of all the **result** expressions must be convertible to a single output type. See Section 10.5 for more details.

There is a “simple” form of **CASE** expression that is a variant of the general form above:

```
CASE expression
    WHEN value THEN result
    [WHEN ...]
    [ELSE result]
END
```

The first **expression** is computed, then compared to each of the **value** expressions in the **WHEN** clauses until one is found that is equal to it. If no match is found, the **result** of the **ELSE** clause (or a null value) is returned. This is similar to the **switch** statement in C.

The example above can be written using the simple **CASE** syntax:
Functions and Operators

SELECT a,
    CASE a WHEN 1 THEN 'one'
        WHEN 2 THEN 'two'
        ELSE 'other'
    END
FROM test;

<table>
<thead>
<tr>
<th>a</th>
<th>case</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
</tr>
<tr>
<td>3</td>
<td>other</td>
</tr>
</tbody>
</table>

A CASE expression does not evaluate any subexpressions that are not needed to determine the result. For example, this is a possible way of avoiding a division-by-zero failure:

SELECT ... WHERE CASE WHEN x <> 0 THEN y/x > 1.5 ELSE false END;

Note

As described in Section 4.2.14, there are various situations in which subexpressions of an expression are evaluated at different times, so that the principle that "CASE evaluates only necessary subexpressions" is not ironclad. For example a constant 1/0 subexpression will usually result in a division-by-zero failure at planning time, even if it's within a CASE arm that would never be entered at run time.

9.17.2. COALESCE

COALESCE(value [, ...])

The COALESCE function returns the first of its arguments that is not null. Null is returned only if all arguments are null. It is often used to substitute a default value for null values when data is retrieved for display, for example:

SELECT COALESCE(description, short_description, '(none)') ...

This returns description if it is not null, otherwise short_description if it is not null, otherwise (none).

Like a CASE expression, COALESCE only evaluates the arguments that are needed to determine the result; that is, arguments to the right of the first non-null argument are not evaluated. This SQL-standard function provides capabilities similar to NVL and IFNULL, which are used in some other database systems.

9.17.3. NULLIF

NULLIF(value1, value2)

The NULLIF function returns a null value if value1 equals value2; otherwise it returns value1. This can be used to perform the inverse operation of the COALESCE example given above:

SELECT NULLIF(value, '(none)') ...
In this example, if `value` is (none), null is returned, otherwise the value of `value` is returned.

### 9.17.4. GREATEST and LEAST

**GREATEST(value [, ...])**

**LEAST(value [, ...])**

The **GREATEST** and **LEAST** functions select the largest or smallest value from a list of any number of expressions. The expressions must all be convertible to a common data type, which will be the type of the result (see Section 10.5 for details). NULL values in the list are ignored. The result will be NULL only if all the expressions evaluate to NULL.

Note that **GREATEST** and **LEAST** are not in the SQL standard, but are a common extension. Some other databases make them return NULL if any argument is NULL, rather than only when all are NULL.

### 9.18. Array Functions and Operators

Table 9.48 shows the operators available for array types.

**Table 9.48. Array Operators**

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>=</td>
<td>equal</td>
<td>ARRAY[1.1,2.1,3.1]::integer = ARRAY[1,2,3]</td>
<td>t</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>not equal</td>
<td>ARRAY[1,2,3] &lt;&gt; ARRAY[1,2,4]</td>
<td>t</td>
</tr>
<tr>
<td>&lt;</td>
<td>less than</td>
<td>ARRAY[1,2,3] &lt; ARRAY[1,2,4]</td>
<td>t</td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
<td>ARRAY[1,4,3] &gt; ARRAY[1,2,4]</td>
<td>t</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal</td>
<td>ARRAY[1,2,3] &lt;= ARRAY[1,2,3]</td>
<td>t</td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater than or equal</td>
<td>ARRAY[1,4,3] &gt;= ARRAY[1,4,3]</td>
<td>t</td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains</td>
<td>ARRAY[1,4,3] @&gt; ARRAY[3,1]</td>
<td>t</td>
</tr>
<tr>
<td>&lt;@</td>
<td>is contained by</td>
<td>ARRAY[2,7] &lt;@ ARRAY[1,7,4,2,6]</td>
<td>t</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>overlap (have elements in common)</td>
<td>ARRAY[1,4,3] &amp;&amp; ARRAY[2,1]</td>
<td>t</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>array-to-array concatenation</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>array-to-array concatenation</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>element-to-array concatenation</td>
</tr>
</tbody>
</table>
Array comparisons compare the array contents element-by-element, using the default B-tree comparison function for the element data type. In multidimensional arrays the elements are visited in row-major order (last subscript varies most rapidly). If the contents of two arrays are equal but the dimensionality is different, the first difference in the dimensionality information determines the sort order. (This is a change from versions of PostgreSQL prior to 8.2: older versions would claim that two arrays with the same contents were equal, even if the number of dimensions or subscript ranges were different.)

See Section 8.15 for more details about array operator behavior. See Section 11.2 for more details about which operators support indexed operations.

Table 9.49 shows the functions available for use with array types. See Section 8.15 for more information and examples of the use of these functions.

**Table 9.49. Array Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>array_append(an_array, anyelement)</td>
<td>anyarray</td>
<td>append an element to the end of an array</td>
<td>array_append(ARRAY[1,2,3], 4)</td>
<td>(1,2,3)</td>
</tr>
<tr>
<td>array_cat(an_array, an_array)</td>
<td>anyarray</td>
<td>concatenate two arrays</td>
<td>array_cat(ARRAY[1,2,3], ARRAY[4,5])</td>
<td>(1,2,3,4,5)</td>
</tr>
<tr>
<td>array_ndims(an_array)</td>
<td>int</td>
<td>returns the number of dimensions of the array</td>
<td>array_ndims(ARRAY[1,2,3], 1)</td>
<td>2</td>
</tr>
<tr>
<td>array_dims(an_array)</td>
<td>text</td>
<td>returns a text representation of array's dimensions</td>
<td>array_dims(ARRAY[1,2,3], 1)</td>
<td>[1:2][1:3]</td>
</tr>
<tr>
<td>array_fill(element, int[], int[])</td>
<td>anyarray</td>
<td>returns an array initialized with supplied value and dimensions, optionally with lower bounds other than 1</td>
<td>array_fill(7, ARRAY[1,2,3], ARRAY[4,5,6])</td>
<td>[2:4]={7,7,7}</td>
</tr>
<tr>
<td>array_length(an_array, int)</td>
<td>int</td>
<td>returns the length of the requested array dimension</td>
<td>array_length(ARRAY[1,2,3], 1)</td>
<td>3</td>
</tr>
<tr>
<td>array_lower(an_array, int)</td>
<td>int</td>
<td>returns lower bound of the requested array dimension</td>
<td>array_lower(ARRAY[1,2,3], 1)</td>
<td>0</td>
</tr>
<tr>
<td>array_position(an_array)</td>
<td>int</td>
<td>returns the subscript of the first occurrence of the second argument 'mon' in the array</td>
<td>array_position(ARRAY['sun','mon','tue','wed','thu','fri'], 'mon')</td>
<td>0</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>----------------------------------------------</td>
<td>--------------</td>
</tr>
<tr>
<td><code>anym element [, int])</code></td>
<td></td>
<td>in the array, starting at the element indicated by the third argument or at the first element (array must be one-dimensional)</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>array_positions(an- yarray, anyelement)</code></td>
<td>int[]</td>
<td>returns an array of subscripts of all occurrences of 'A', 'A', 'B', 'A' (the second argument in the array given as first argument (array must be one-dimensional))</td>
<td>(1,2,4)</td>
<td></td>
</tr>
<tr>
<td><code>ar- ray_prepend(an- yarray)</code></td>
<td>anyarray</td>
<td>append an element to the beginning of an array (array must be one-dimensional)</td>
<td>(1,2,3)</td>
<td></td>
</tr>
<tr>
<td><code>array_re- move(an- yarray, anyelement)</code></td>
<td>anyarray</td>
<td>remove all elements equal to the given value from the array (array must be one-dimensional)</td>
<td>(1,3)</td>
<td></td>
</tr>
<tr>
<td><code>array_re- place(an- yarray, anyelement, anyelement)</code></td>
<td>anyarray</td>
<td>replace each array element equal to the given value with a new value</td>
<td>(1,2,3,4)</td>
<td></td>
</tr>
<tr>
<td><code>ar- ray_to_string(an- yarray, text [, text])</code></td>
<td>text</td>
<td>concatenates array elements using supplied delimiter and optional null string</td>
<td>(1,2,3,*,5)</td>
<td></td>
</tr>
<tr>
<td><code>array_up- per(an- yarray, int)</code></td>
<td>int</td>
<td>returns the upper bound of the requested array dimension</td>
<td>(1,8,3,7), 1</td>
<td></td>
</tr>
<tr>
<td><code>cardinali- ty(anyarray)</code></td>
<td>int</td>
<td>returns the total number of elements in the array.</td>
<td>(1,2,3), [3,4])</td>
<td></td>
</tr>
<tr>
<td><code>string_to_ar- ray(text, text [, text])</code></td>
<td></td>
<td>splits string into array elements using supplied delimiter and optional null string</td>
<td>('x','y','z'), 'yy'</td>
<td></td>
</tr>
<tr>
<td><code>unnest(an- yarray)</code></td>
<td>setof anyele- ment</td>
<td>expand an array to a set of rows</td>
<td>(AR-[1,2])</td>
<td></td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>unnest(anyarray, anyarray) [, ...])</td>
<td>setof anyelement, anyelement [, ...]</td>
<td>expand multiple arrays (possibly of different types) to a set of rows. This is only allowed in the FROM clause; see Section 7.2.1.4</td>
<td>unnest(ARRAY[1,2],ARRAY['foo','bar','baz'])</td>
<td>1 foo 2 bar NULL baz (3 rows)</td>
</tr>
</tbody>
</table>

In `array_position` and `array_positions`, each array element is compared to the searched value using IS NOT DISTINCT FROM semantics.

In `array_position`, NULL is returned if the value is not found.

In `array_positions`, NULL is returned only if the array is NULL; if the value is not found in the array, an empty array is returned instead.

In `string_to_array`, if the delimiter parameter is NULL, each character in the input string will become a separate element in the resulting array. If the delimiter is an empty string, then the entire input string is returned as a one-element array. Otherwise the input string is split at each occurrence of the delimiter string.

In `string_to_array`, if the null-string parameter is omitted or NULL, none of the substrings of the input will be replaced by NULL. In `array_to_string`, if the null-string parameter is omitted or NULL, any null elements in the array are simply skipped and not represented in the output string.

#### Note

There are two differences in the behavior of `string_to_array` from pre-9.1 versions of PostgreSQL. First, it will return an empty (zero-element) array rather than NULL when the input string is of zero length. Second, if the delimiter string is NULL, the function splits the input into individual characters, rather than returning NULL as before.

See also Section 9.20 about the aggregate function `array_agg` for use with arrays.

### 9.19. Range Functions and Operators

See Section 8.17 for an overview of range types.

Table 9.50 shows the operators available for range types.

#### Table 9.50. Range Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>=</td>
<td>equal</td>
<td>int4range(1,5) = t '1,4':'::int4range</td>
<td>t</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>not equal</td>
<td>numrange(1.1,2.2) &lt;&gt; numrange(1.1,2.3)</td>
<td>t</td>
</tr>
<tr>
<td>&lt;</td>
<td>less than</td>
<td>int4range(1,10) &lt; int4range(2,3)</td>
<td>t</td>
</tr>
<tr>
<td>Operator</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>------------------------------</td>
<td>------------------</td>
<td>--------</td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
<td><code>int4range(1,10) &gt; int4range(1,5)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal</td>
<td><code>numrange(1.1,2.2) &lt;= numrange(1.1,2.2)</code></td>
<td>t</td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater than or equal</td>
<td><code>numrange(1.1,2.2) &gt;= numrange(1.1,2.0)</code></td>
<td>t</td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains range</td>
<td><code>int4range(2,4) @&gt; in- t4range(2,3)</code></td>
<td>t</td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains element</td>
<td>'[2011-01-01,2011-03-01)::tsrange @&gt; '2011-01-10':::time-stamp'</td>
<td>t</td>
</tr>
<tr>
<td>&lt;@</td>
<td>range is contained by</td>
<td><code>int4range(2,4) t &lt;@ in- t4range(1,7)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;@</td>
<td>element is contained by</td>
<td>42 &lt;@ in-f t4range(1,7)</td>
<td>t</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>overlap (have points in common)</td>
<td><code>int8range(3,7) &amp;&amp; in- t8range(4,12)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>strictly left of</td>
<td><code>int8range(1,10) t &lt;&lt; in- t8range(100,110)</code></td>
<td>t</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>strictly right of</td>
<td><code>int8range(50,60) t &gt;&gt; in- t8range(20,30)</code></td>
<td>t</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td>does not extend to the right of</td>
<td><code>int8range(1,20) t &amp;&lt; in- t8range(18,20)</code></td>
<td>t</td>
</tr>
<tr>
<td>&amp;&gt;</td>
<td>does not extend to the left of</td>
<td><code>int8range(7,20) t &amp;&gt; in- t8range(5,10)</code></td>
<td>t</td>
</tr>
<tr>
<td>-=</td>
<td>is adjacent to</td>
<td><code>numrange(1.1,2.2) -= numrange(2.2,3.3)</code></td>
<td>t</td>
</tr>
<tr>
<td>+</td>
<td>union</td>
<td><code>numrange(5,15) + [5,20]</code></td>
<td><code>numrange(10,20)</code></td>
</tr>
<tr>
<td>*</td>
<td>intersection</td>
<td><code>int8range(5,15) [10,15) * in- t8range(10,20)</code></td>
<td>t</td>
</tr>
<tr>
<td>-</td>
<td>difference</td>
<td><code>int8range(5,15) [5,10) - in- t8range(10,20)</code></td>
<td>t</td>
</tr>
</tbody>
</table>
The simple comparison operators `<`, `>`, `<=`, and `>=` compare the lower bounds first, and only if those are equal, compare the upper bounds. These comparisons are not usually very useful for ranges, but are provided to allow B-tree indexes to be constructed on ranges.

The left-of/right-of/adjacent operators always return false when an empty range is involved; that is, an empty range is not considered to be either before or after any other range.

The union and difference operators will fail if the resulting range would need to contain two disjoint sub-ranges, as such a range cannot be represented.

Table 9.51 shows the functions available for use with range types.

### Table 9.51. Range Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>lower(anyrange)</code></td>
<td>range's element type</td>
<td>lower bound of range</td>
<td><code>lower(num_range(1.1,2.2))</code></td>
<td>1.1</td>
</tr>
<tr>
<td><code>upper(anyrange)</code></td>
<td>range's element type</td>
<td>upper bound of range</td>
<td><code>upper(num_range(1.1,2.2))</code></td>
<td>2.2</td>
</tr>
<tr>
<td><code>isempty(anyrange)</code></td>
<td>boolean</td>
<td>is the range empty?</td>
<td><code>isempty(num_range(1.1,2.2))</code></td>
<td>false</td>
</tr>
<tr>
<td><code>lower_inc(anyrange)</code></td>
<td>boolean</td>
<td>is the lower bound inclusive?</td>
<td><code>lower_inc(num_range(1.1,2.2))</code></td>
<td>true</td>
</tr>
<tr>
<td><code>upper_inc(anyrange)</code></td>
<td>boolean</td>
<td>is the upper bound inclusive?</td>
<td><code>upper_inc(num_range(1.1,2.2))</code></td>
<td>false</td>
</tr>
<tr>
<td><code>lower_inf(anyrange)</code></td>
<td>boolean</td>
<td>is the lower bound infinite?</td>
<td><code>lower_inf('(,)'::daterange)</code></td>
<td>true</td>
</tr>
<tr>
<td><code>upper_inf(anyrange)</code></td>
<td>boolean</td>
<td>is the upper bound infinite?</td>
<td><code>upper_inf('(,)'::daterange)</code></td>
<td>true</td>
</tr>
<tr>
<td><code>range_merge(anyrange, anyrange)</code></td>
<td>anyrange</td>
<td>the smallest range which includes both of the given ranges</td>
<td><code>range_merge('([1,2])'::int4range, '[3,4]'::int4range)</code></td>
<td></td>
</tr>
</tbody>
</table>

The `lower` and `upper` functions return null if the range is empty or the requested bound is infinite. The `lower_inc`, `upper_inc`, `lower_inf`, and `upper_inf` functions all return false for an empty range.

### 9.20. Aggregate Functions

*Aggregate functions* compute a single result from a set of input values. The built-in general-purpose aggregate functions are listed in Table 9.52 and statistical aggregates in Table 9.53. The built-in within-group ordered-set aggregate functions are listed in Table 9.54 while the built-in within-group hypothetical-set ones are in Table 9.55. Grouping operations, which are closely related to aggregate functions, are listed in Table 9.56. The special syntax considerations for aggregate functions are explained in Section 4.2.7. Consult Section 2.7 for additional introductory information.
### Table 9.52. General-Purpose Aggregate Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>array_agg(expression)</code></td>
<td>any non-array type</td>
<td>array of the argument type</td>
<td>No</td>
<td>input values, including nulls, concatenated into an array</td>
</tr>
<tr>
<td><code>array_agg(expression)</code></td>
<td>any array type</td>
<td>same as argument data type</td>
<td>No</td>
<td>input arrays concatenated into an array of one higher dimension (inputs must all have the same dimensionality, and cannot be empty or NULL)</td>
</tr>
<tr>
<td><code>avg(expression)</code></td>
<td>smallint, int, numeric for any bigint, real, integer-type argument, double precision, numeric, precision for a floating-point argument, otherwise the same as the argument data type</td>
<td>numeric for any bigint, real, integer-type argument, double precision, numeric, precision for a floating-point argument, otherwise the same as the argument data type</td>
<td>Yes</td>
<td>the average (arithmetic mean) of all input values</td>
</tr>
<tr>
<td><code>bit_and(expression)</code></td>
<td>smallint, int, same as argument data type</td>
<td>bigint, or bit</td>
<td>Yes</td>
<td>the bitwise AND of all non-null input values, or null if none</td>
</tr>
<tr>
<td><code>bit_or(expression)</code></td>
<td>smallint, int, same as argument data type</td>
<td>bigint, or bit</td>
<td>Yes</td>
<td>the bitwise OR of all non-null input values, or null if none</td>
</tr>
<tr>
<td><code>bool_and(expression)</code></td>
<td>bool</td>
<td>bool</td>
<td>Yes</td>
<td>true if all input values are true, otherwise false</td>
</tr>
<tr>
<td><code>bool_or(expression)</code></td>
<td>bool</td>
<td>bool</td>
<td>Yes</td>
<td>true if at least one input value is true, otherwise false</td>
</tr>
<tr>
<td><code>count(*)</code></td>
<td></td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows</td>
</tr>
<tr>
<td><code>count(expression)</code></td>
<td>any</td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows for which the value of expression is not null</td>
</tr>
<tr>
<td><code>every(expression)</code></td>
<td></td>
<td>bool</td>
<td>Yes</td>
<td>equivalent to <code>bool_and</code></td>
</tr>
<tr>
<td><code>json_agg(expression)</code></td>
<td>any</td>
<td>json</td>
<td>No</td>
<td>aggregates values as a JSON array</td>
</tr>
<tr>
<td><code>jsonb_agg(expression)</code></td>
<td>any</td>
<td>jsonb</td>
<td>No</td>
<td>aggregates values as a JSON array</td>
</tr>
<tr>
<td>Function</td>
<td>Argument Type(s)</td>
<td>Return Type</td>
<td>Partial Mode</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>------------------</td>
<td>-------------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>json_object_agg</td>
<td>(any, any)</td>
<td>json</td>
<td>No</td>
<td>aggregates name/value pairs as a JSON object</td>
</tr>
<tr>
<td>jsonb_object_agg</td>
<td>(any, any)</td>
<td>jsonb</td>
<td>No</td>
<td>aggregates name/value pairs as a JSON object</td>
</tr>
<tr>
<td>max(expression)</td>
<td>any numeric, same as argument</td>
<td>same as argument</td>
<td>Yes</td>
<td>maximum value of expression across all input values</td>
</tr>
<tr>
<td>min(expression)</td>
<td>any numeric, same as argument</td>
<td>same as argument</td>
<td>Yes</td>
<td>minimum value of expression across all input values</td>
</tr>
<tr>
<td>string_agg</td>
<td>(text, text) or same as argument (bytea, bytea) types</td>
<td>same as argument</td>
<td>No</td>
<td>input values concatenated into a string, separated by delimiter</td>
</tr>
<tr>
<td>sum(expression)</td>
<td>smallint, int, bigint or double precision int arguments, numeric for ic, interval, bigint arguments, otherwise the same as the argument data type</td>
<td>bigint for smallint or int arguments, numeric for bigint arguments, otherwise the same as the argument data type</td>
<td>Yes</td>
<td>sum of expression across all input values</td>
</tr>
<tr>
<td>xmlagg(expression)</td>
<td>xml</td>
<td>No</td>
<td>Concatenation of XML values (see also Section 9.14.1.7)</td>
<td></td>
</tr>
</tbody>
</table>

It should be noted that except for `count`, these functions return a null value when no rows are selected. In particular, `sum` of no rows returns null, not zero as one might expect, and `array_agg` returns null rather than an empty array when there are no input rows. The `coalesce` function can be used to substitute zero or an empty array for null when necessary.

Aggregate functions which support `Partial Mode` are eligible to participate in various optimizations, such as parallel aggregation.

**Note**

Boolean aggregates `bool_and` and `bool_or` correspond to standard SQL aggregates `every` and `any` or `some`. As for `any` and `some`, it seems that there is an ambiguity built into the standard syntax:

```
SELECT b1 = ANY((SELECT b2 FROM t2 ...)) FROM t1 ...;
```
Here \( \text{ANY} \) can be considered either as introducing a subquery, or as being an aggregate function, if the subquery returns one row with a Boolean value. Thus the standard name cannot be given to these aggregates.

**Note**

Users accustomed to working with other SQL database management systems might be disappointed by the performance of the `count` aggregate when it is applied to the entire table. A query like:

```
SELECT count(*) FROM sometable;
```

will require effort proportional to the size of the table: PostgreSQL will need to scan either the entire table or the entirety of an index which includes all rows in the table.

The aggregate functions `array_agg`, `json_agg`, `jsonb_agg`, `json_object_agg`, `jsonb_object_agg`, `string_agg`, and `xmlagg`, as well as similar user-defined aggregate functions, produce meaningfully different result values depending on the order of the input values. This ordering is unspecified by default, but can be controlled by writing an `ORDER BY` clause within the aggregate call, as shown in Section 4.2.7. Alternatively, supplying the input values from a sorted subquery will usually work. For example:

```
SELECT xmlagg(x) FROM (SELECT x FROM test ORDER BY y DESC) AS tab;
```

Beware that this approach can fail if the outer query level contains additional processing, such as a join, because that might cause the subquery’s output to be reordered before the aggregate is computed.

Table 9.53 shows aggregate functions typically used in statistical analysis. (These are separated out merely to avoid cluttering the listing of more-commonly-used aggregates.) Where the description mentions \( N \), it means the number of input rows for which all the input expressions are non-null. In all cases, null is returned if the computation is meaningless, for example when \( N \) is zero.

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>corr(Y,X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>correlation coefficient</td>
</tr>
<tr>
<td><code>covar_pop(Y,X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>population covariance</td>
</tr>
<tr>
<td><code>covar_samp(Y,X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>sample covariance</td>
</tr>
<tr>
<td><code>regr_avgx(Y,X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>average of the independent variable ((\text{sum}(X) / N))</td>
</tr>
<tr>
<td><code>regr_avgy(Y,X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>average of the dependent variable ((\text{sum}(Y) / N))</td>
</tr>
<tr>
<td><code>regr_count(Y,X)</code></td>
<td>double precision</td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows in which both expressions are nonnull</td>
</tr>
<tr>
<td>Function</td>
<td>Argument Type</td>
<td>Return Type</td>
<td>Partial Mode</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>-----------------------</td>
<td>---------------------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>regr_intercept(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>y-intercept of the least-squares-fit linear equation determined by the (X, Y) pairs</td>
</tr>
<tr>
<td><code>regr_r2(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>square of the correlation coefficient</td>
</tr>
<tr>
<td><code>regr_slope(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>slope of the least-squares-fit linear equation determined by the (X, Y) pairs</td>
</tr>
<tr>
<td><code>regr_sxx(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td><code>sum(X^2) - sum(X)^2/N</code> (&quot;sum of squares&quot; of the independent variable)</td>
</tr>
<tr>
<td><code>regr_sxy(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td><code>sum(X*Y) - sum(X)*sum(Y)/N</code> (&quot;sum of products&quot; of independent times dependent variable)</td>
</tr>
<tr>
<td><code>regr_syy(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td><code>sum(Y^2) - sum(Y)^2/N</code> (&quot;sum of squares&quot; of the dependent variable)</td>
</tr>
<tr>
<td><code>stddev(expression)</code></td>
<td>smallint, int, double precision, bigint, real, cision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>historical alias for stddev_samp</td>
</tr>
<tr>
<td><code>var_pop(expression)</code></td>
<td>smallint, int, double precision, bigint, real, cision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>population standard deviation of the input values</td>
</tr>
<tr>
<td><code>var_samp(expression)</code></td>
<td>smallint, int, double precision, bigint, real, cision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>sample standard deviation of the input values</td>
</tr>
<tr>
<td><code>variance(expression)</code></td>
<td>smallint, int, double precision, bigint, real, cision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>historical alias for var_samp</td>
</tr>
<tr>
<td><code>var_pop(expression)</code></td>
<td>smallint, int, double precision, bigint, real, cision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>population variance of the input values (square root of <code>var_pop</code>)</td>
</tr>
<tr>
<td>Function</td>
<td>Argument Type</td>
<td>Return Type</td>
<td>Partial Mode</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>---------------</td>
<td>-------------</td>
<td>--------------</td>
<td>-------------</td>
</tr>
<tr>
<td>var_samp(expression)</td>
<td>smallint, int, double precision for floating-point arguments, otherwise numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>sample variance of the input values (square of the sample standard deviation)</td>
</tr>
</tbody>
</table>

Table 9.54 shows some aggregate functions that use the ordered-set aggregate syntax. These functions are sometimes referred to as “inverse distribution” functions.

### Table 9.54. Ordered-Set Aggregate Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Direct Argument Type(s)</th>
<th>Aggregated Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>mode()</td>
<td>WITHIN GROUP (ORDER BY sort_expression)</td>
<td>any sortable type</td>
<td>same as sort expression</td>
<td>No</td>
<td>returns the most frequent input value (arbitrarily choosing the first one if there are multiple equally-frequent results)</td>
</tr>
<tr>
<td>percentile_cont(fraction)</td>
<td>double precision</td>
<td>double precision or interval</td>
<td>same as sort expression</td>
<td>No</td>
<td>continuous percentile: returns a value corresponding to the specified fraction in the ordering, interpolating between adjacent input items if needed</td>
</tr>
<tr>
<td>percentile_cont(fractions)</td>
<td>double precision[]</td>
<td>double precision or interval</td>
<td>array of sort expression's type</td>
<td>No</td>
<td>multiple continuous percentile: returns an array of results matching the shape of the fractions parameter, with each non-null element replaced by the value corresponding to that percentile</td>
</tr>
<tr>
<td>percentile_disc(fraction)</td>
<td>double precision</td>
<td>any sortable type</td>
<td>same as sort expression</td>
<td>No</td>
<td>discrete percentile: returns the first input value whose...</td>
</tr>
</tbody>
</table>
Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Direct Argument Type(s)</th>
<th>Aggregated Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>WITHIN GROUP (ORDER BY sort_expression)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>position in the ordering equals or exceeds the specified fraction</td>
</tr>
<tr>
<td>percentile_distributions</td>
<td>double precision</td>
<td>any sortable type</td>
<td>array of sort expression's type</td>
<td>No</td>
<td>multiple discrete percentile: returns an array of results matching the shape of the fractions parameter, with each non-null element replaced by the input value corresponding to that percentile</td>
</tr>
</tbody>
</table>

All the aggregates listed in Table 9.54 ignore null values in their sorted input. For those that take a fraction parameter, the fraction value must be between 0 and 1; an error is thrown if not. However, a null fraction value simply produces a null result.

Each of the aggregates listed in Table 9.55 is associated with a window function of the same name defined in Section 9.21. In each case, the aggregate result is the value that the associated window function would have returned for the “hypothetical” row constructed from args, if such a row had been added to the sorted group of rows computed from the sorted_args.

Table 9.55. Hypothetical-Set Aggregate Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Direct Argument Type(s)</th>
<th>Aggregated Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rank(args) WITHIN GROUP (ORDER BY sorted_args)</td>
<td>VARIADIC &quot;any&quot;</td>
<td>VARIADIC &quot;any&quot;</td>
<td>bigint</td>
<td>No</td>
<td>rank of the hypothetical row, with gaps for duplicate rows</td>
</tr>
<tr>
<td>dense_rank(args) WITHIN GROUP (ORDER BY sorted_args)</td>
<td>VARIADIC &quot;any&quot;</td>
<td>VARIADIC &quot;any&quot;</td>
<td>bigint</td>
<td>No</td>
<td>rank of the hypothetical row, without gaps</td>
</tr>
<tr>
<td>percent_rank(args) WITHIN GROUP (ORDER BY</td>
<td>VARIADIC &quot;any&quot;</td>
<td>VARIADIC &quot;any&quot;</td>
<td>double precision</td>
<td>No</td>
<td>relative rank of the hypothetical row, ranging from 0 to 1</td>
</tr>
</tbody>
</table>
For each of these hypothetical-set aggregates, the list of direct arguments given in \( args \) must match the number and types of the aggregated arguments given in \( \text{sorted}_\text{args} \). Unlike most built-in aggregates, these aggregates are not strict, that is they do not drop input rows containing nulls. Null values sort according to the rule specified in the \( \text{ORDER BY} \) clause.

### Table 9.56. Grouping Operations

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>( \text{GROUPING}(\text{args}...) )</td>
<td>integer</td>
<td>integer bit mask indicating which arguments are not being included in the current grouping set</td>
</tr>
</tbody>
</table>

Grouping operations are used in conjunction with grouping sets (see Section 7.2.4) to distinguish result rows. The arguments to the \( \text{GROUPING} \) operation are not actually evaluated, but they must match exactly expressions given in the \( \text{GROUP BY} \) clause of the associated query level. Bits are assigned with the rightmost argument being the least-significant bit; each bit is 0 if the corresponding expression is included in the grouping criteria of the grouping set generating the result row, and 1 if it is not. For example:

```sql
=> SELECT * FROM items_sold;
make | model | sales
-------+-------+-------
Foo   | GT    | 10    
Foo   | Tour  | 20    
Bar   | City  | 15    
Bar   | Sport | 5     
(4 rows)
```

```sql
=> SELECT make, model, \( \text{GROUPING}(\text{make},\text{model}) \), \( \text{sum}(\text{sales}) \) FROM items_sold GROUP BY ROLLUP(make, model);
make | model | grouping | sum
-------+-------+----------+-----
Foo   | GT    | 0        | 10  
Foo   | Tour  | 0        | 20  
Bar   | City  | 0        | 15  
Bar   | Sport | 0        | 5   
Foo   |       | 1        | 30  
Bar   |       | 1        | 20  
|       |        | 3        | 50  
(7 rows)
```

### 9.21. Window Functions
Window functions provide the ability to perform calculations across sets of rows that are related to the current query row. See Section 3.5 for an introduction to this feature, and Section 4.2.8 for syntax details.

The built-in window functions are listed in Table 9.57. Note that these functions must be invoked using window function syntax, i.e., an OVER clause is required.

In addition to these functions, any built-in or user-defined general-purpose or statistical aggregate (i.e., not ordered-set or hypothetical-set aggregates) can be used as a window function; see Section 9.20 for a list of the built-in aggregates. Aggregate functions act as window functions only when an OVER clause follows the call; otherwise they act as non-window aggregates and return a single row for the entire set.

<table>
<thead>
<tr>
<th>Table 9.57. General-Purpose Window Functions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Function</td>
</tr>
<tr>
<td>row_number()</td>
</tr>
<tr>
<td>rank()</td>
</tr>
<tr>
<td>dense_rank()</td>
</tr>
<tr>
<td>percent_rank()</td>
</tr>
<tr>
<td>cume_dist()</td>
</tr>
<tr>
<td>ntile(num_buckets integer)</td>
</tr>
<tr>
<td>lag(value anyelement [, offset integer [, default anyelement ]])</td>
</tr>
<tr>
<td>lead(value anyelement [, offset integer [, default anyelement ]])</td>
</tr>
</tbody>
</table>
Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>first_value(value any)</td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the first row of the window frame</td>
</tr>
<tr>
<td>last_value(value any)</td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the last row of the window frame</td>
</tr>
<tr>
<td>nth_value(value any, nth integer)</td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the nth row of the window frame (counting from 1); null if no such row</td>
</tr>
</tbody>
</table>

All of the functions listed in Table 9.57 depend on the sort ordering specified by the ORDER BY clause of the associated window definition. Rows that are not distinct when considering only the ORDER BY columns are said to be peers. The four ranking functions (including cume_dist) are defined so that they give the same answer for all peer rows.

Note that first_value, last_value, and nth_value consider only the rows within the “window frame”, which by default contains the rows from the start of the partition through the last peer of the current row. This is likely to give unhelpful results for last_value and sometimes also nth_value. You can redefine the frame by adding a suitable frame specification (RANGE, ROWS or GROUPS) to the OVER clause. See Section 4.2.8 for more information about frame specifications.

When an aggregate function is used as a window function, it aggregates over the rows within the current row’s window frame. An aggregate used with ORDER BY and the default window frame definition produces a “running sum” type of behavior, which may or may not be what’s wanted. To obtain aggregation over the whole partition, omit ORDER BY or use ROWS BETWEEN UNBOUNDED PRECEDING AND UNBOUNDED FOLLOWING. Other frame specifications can be used to obtain other effects.

Note

The SQL standard defines a RESPECT NULLS or IGNORE NULLS option for lead, lag, first_value, last_value, and nth_value. This is not implemented in PostgreSQL: the behavior is always the same as the standard’s default, namely RESPECT NULLS. Likewise, the standard’s FROM FIRST or FROM LAST option for nth_value is not implemented: only the default FROM FIRST behavior is supported. (You can achieve the result of FROM LAST by reversing the ORDER BY ordering.)

cume_dist computes the fraction of partition rows that are less than or equal to the current row and its peers, while percent_rank computes the fraction of partition rows that are less than the current row, assuming the current row does not exist in the partition.

9.22. Subquery Expressions

This section describes the SQL-compliant subquery expressions available in PostgreSQL. All of the expression forms documented in this section return Boolean (true/false) results.

9.22.1. EXISTS

EXISTS (subquery)

The argument of EXISTS is an arbitrary SELECT statement, or subquery. The subquery is evaluated to determine whether it returns any rows. If it returns at least one row, the result of EXISTS is “true”; if the subquery returns no rows, the result of EXISTS is “false”.

312
The subquery can refer to variables from the surrounding query, which will act as constants during any one evaluation of the subquery.

The subquery will generally only be executed long enough to determine whether at least one row is returned, not all the way to completion. It is unwise to write a subquery that has side effects (such as calling sequence functions); whether the side effects occur might be unpredictable.

Since the result depends only on whether any rows are returned, and not on the contents of those rows, the output list of the subquery is normally unimportant. A common coding convention is to write all EXISTS tests in the form EXISTS(SELECT 1 WHERE ...). There are exceptions to this rule however, such as subqueries that use INTERSECT.

This simple example is like an inner join on col2, but it produces at most one output row for each tab1 row, even if there are several matching tab2 rows:

```
SELECT col1
FROM tab1
WHERE EXISTS (SELECT 1 FROM tab2 WHERE col2 = tab1.col2);
```

### 9.22.2. IN

expression IN (subquery)

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result. The result of IN is “true” if any equal subquery row is found. The result is “false” if no equal row is found (including the case where the subquery returns no rows).

Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand row yields null, the result of the IN construct will be null, not false. This is in accordance with SQL's normal rules for Boolean combinations of null values.

As with EXISTS, it’s unwise to assume that the subquery will be evaluated completely.

row_constructor IN (subquery)

The left-hand side of this form of IN is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result. The result of IN is “true” if any equal subquery row is found. The result is “false” if no equal row is found (including the case where the subquery returns no rows).

As usual, null values in the rows are combined per the normal rules of SQL Boolean expressions. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of that row comparison is unknown (null). If all the per-row results are either unequal or null, with at least one null, then the result of IN is null.

### 9.22.3. NOT IN

expression NOT IN (subquery)

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result. The result of NOT IN is “true” if only unequal subquery rows are found (including the case where the subquery returns no rows). The result is “false” if any equal row is found.
Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand row yields null, the result of the \texttt{NOT IN} construct will be null, not true. This is in accordance with SQL's normal rules for Boolean combinations of null values.

As with \texttt{EXISTS}, it's unwise to assume that the subquery will be evaluated completely.

\texttt{row\_constructor \texttt{NOT IN} (subquery)}

The left-hand side of this form of \texttt{NOT IN} is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result. The result of \texttt{NOT IN} is “true” if only unequal subquery rows are found (including the case where the subquery returns no rows). The result is “false” if any equal row is found.

As usual, null values in the rows are combined per the normal rules of SQL Boolean expressions. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of that row comparison is unknown (null). If all the per-row results are either unequal or null, with at least one null, then the result of \texttt{NOT IN} is null.

\subsection{ANY/SOME}

\texttt{expression \texttt{operator ANY} (subquery)}
\texttt{expression \texttt{operator SOME} (subquery)}

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result using the given \texttt{operator}, which must yield a Boolean result. The result of \texttt{ANY} is “true” if any true result is obtained. The result is “false” if no true result is found (including the case where the subquery returns no rows).

\texttt{SOME} is a synonym for \texttt{ANY}. \texttt{IN} is equivalent to \texttt{= ANY}.

Note that if there are no successes and at least one right-hand row yields null for the operator's result, the result of the \texttt{ANY} construct will be null, not false. This is in accordance with SQL's normal rules for Boolean combinations of null values.

As with \texttt{EXISTS}, it's unwise to assume that the subquery will be evaluated completely.

\texttt{row\_constructor \texttt{operator ANY} (subquery)}
\texttt{row\_constructor \texttt{operator SOME} (subquery)}

The left-hand side of this form of \texttt{ANY} is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result, using the given \texttt{operator}. The result of \texttt{ANY} is “true” if the comparison returns true for any subquery row. The result is “false” if the comparison returns false for every subquery row (including the case where the subquery returns no rows). The result is NULL if no comparison with a subquery row returns true, and at least one comparison returns NULL.

See Section 9.23.5 for details about the meaning of a row constructor comparison.

\subsection{ALL}

\texttt{expression \texttt{operator ALL} (subquery)}
The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result using the given operator, which must yield a Boolean result. The result of ALL is "true" if all rows yield true (including the case where the subquery returns no rows). The result is "false" if any false result is found. The result is NULL if no comparison with a subquery row returns false, and at least one comparison returns NULL.

NOT IN is equivalent to <> ALL.

As with EXISTS, it's unwise to assume that the subquery will be evaluated completely.

row_constructor operator ALL (subquery)

The left-hand side of this form of ALL is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result, using the given operator. The result of ALL is "true" if the comparison returns true for all subquery rows (including the case where the subquery returns no rows). The result is "false" if the comparison returns false for any subquery row. The result is NULL if no comparison with a subquery row returns false, and at least one comparison returns NULL.

See Section 9.23.5 for details about the meaning of a row constructor comparison.

9.22.6. Single-row Comparison

row_constructor operator (subquery)

The left-hand side is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. Furthermore, the subquery cannot return more than one row. (If it returns zero rows, the result is taken to be null.) The left-hand side is evaluated and compared row-wise to the single subquery result row.

See Section 9.23.5 for details about the meaning of a row constructor comparison.

9.23. Row and Array Comparisons

This section describes several specialized constructs for making multiple comparisons between groups of values. These forms are syntactically related to the subquery forms of the previous section, but do not involve subqueries. The forms involving array subexpressions are PostgreSQL extensions; the rest are SQL-compliant. All of the expression forms documented in this section return Boolean (true/false) results.

9.23.1. IN

expression IN (value [, ...])

The right-hand side is a parenthesized list of scalar expressions. The result is "true" if the left-hand expression's result is equal to any of the right-hand expressions. This is a shorthand notation for

expression = value1
OR
expression = value2
OR
...


Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand expression yields null, the result of the \texttt{IN} construct will be null, not false. This is in accordance with SQL's normal rules for Boolean combinations of null values.

### 9.23.2. \texttt{NOT IN}

\texttt{expression NOT IN (value [,, ...])}

The right-hand side is a parenthesized list of scalar expressions. The result is “true” if the left-hand expression's result is unequal to all of the right-hand expressions. This is a shorthand notation for

\texttt{expression <> value1}
\texttt{AND}
\texttt{expression <> value2}
\texttt{AND}
\texttt{...}

Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand expression yields null, the result of the \texttt{NOT IN} construct will be null, not true as one might naively expect. This is in accordance with SQL's normal rules for Boolean combinations of null values.

**Tip**

\texttt{x NOT IN y} is equivalent to \texttt{NOT (x IN y)} in all cases. However, null values are much more likely to trip up the novice when working with \texttt{NOT IN} than when working with \texttt{IN}. It is best to express your condition positively if possible.

### 9.23.3. \texttt{ANY/SOME} (array)

\texttt{expression operator ANY (array expression)}
\texttt{expression operator SOME (array expression)}

The right-hand side is a parenthesized expression, which must yield an array value. The left-hand expression is evaluated and compared to each element of the array using the given \texttt{operator}, which must yield a Boolean result. The result of \texttt{ANY} is “true” if any true result is obtained. The result is “false” if no true result is found (including the case where the array has zero elements).

If the array expression yields a null array, the result of \texttt{ANY} will be null. If the left-hand expression yields null, the result of \texttt{ANY} is ordinarily null (though a non-strict comparison operator could possibly yield a different result). Also, if the right-hand array contains any null elements and no true comparison result is obtained, the result of \texttt{ANY} will be null, not false (again, assuming a strict comparison operator). This is in accordance with SQL's normal rules for Boolean combinations of null values.

\texttt{SOME} is a synonym for \texttt{ANY}.

### 9.23.4. \texttt{ALL} (array)

\texttt{expression operator ALL (array expression)}

The right-hand side is a parenthesized expression, which must yield an array value. The left-hand expression is evaluated and compared to each element of the array using the given \texttt{operator}, which must yield a Boolean result. The result of \texttt{ALL} is “true” if all comparisons yield true (including the case where the array has zero elements). The result is “false” if any false result is found.
If the array expression yields a null array, the result of `ALL` will be null. If the left-hand expression yields null, the result of `ALL` is ordinarily null (though a non-strict comparison operator could possibly yield a different result). Also, if the right-hand array contains any null elements and no false comparison result is obtained, the result of `ALL` will be null, not true (again, assuming a strict comparison operator). This is in accordance with SQL’s normal rules for Boolean combinations of null values.

9.23.5. Row Constructor Comparison

`row_constructor operator row_constructor`

Each side is a row constructor, as described in Section 4.2.13. The two row values must have the same number of fields. Each side is evaluated and they are compared row-wise. Row constructor comparisons are allowed when the `operator` is `=`, `<>`, `<=`, `>`, or `>=`. Every row element must be of a type which has a default B-tree operator class or the attempted comparison may generate an error.

Note

Errors related to the number or types of elements might not occur if the comparison is resolved using earlier columns.

The `=` and `<>` cases work slightly differently from the others. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of the row comparison is unknown (null).

For the `<`, `<=`, `>` and `>=` cases, the row elements are compared left-to-right, stopping as soon as an unequal or null pair of elements is found. If either of this pair of elements is null, the result of the row comparison is unknown (null); otherwise comparison of this pair of elements determines the result. For example, `ROW(1,2,NULL) < ROW(1,3,0)` yields true, not null, because the third pair of elements are not considered.

Note

Prior to PostgreSQL 8.2, the `<`, `<=`, `>` and `>=` cases were not handled per SQL specification. A comparison like `ROW(a,b) < ROW(c,d)` was implemented as `a < c AND b < d` whereas the correct behavior is equivalent to `a < c OR (a = c AND b < d)`.

`row_constructor IS DISTINCT FROM row_constructor`

This construct is similar to a `<>` row comparison, but it does not yield null for null inputs. Instead, any null value is considered unequal to (distinct from) any non-null value, and any two nulls are considered equal (not distinct). Thus the result will either be true or false, never null.

`row_constructor IS NOT DISTINCT FROM row_constructor`

This construct is similar to `a =` row comparison, but it does not yield null for null inputs. Instead, any null value is considered unequal to (distinct from) any non-null value, and any two nulls are considered equal (not distinct). Thus the result will always be either true or false, never null.

9.23.6. Composite Type Comparison

`record operator record`
The SQL specification requires row-wise comparison to return NULL if the result depends on comparing two NULL values or a NULL and a non-NULL. PostgreSQL does this only when comparing the results of two row constructors (as in Section 9.23.5) or comparing a row constructor to the output of a subquery (as in Section 9.22). In other contexts where two composite-type values are compared, two NULL field values are considered equal, and a NULL is considered larger than a non-NULL. This is necessary in order to have consistent sorting and indexing behavior for composite types.

Each side is evaluated and they are compared row-wise. Composite type comparisons are allowed when the operator is =, <>, <, <=, or >=, or has semantics similar to one of these. (To be specific, an operator can be a row comparison operator if it is a member of a B-tree operator class, or is the negator of the = member of a B-tree operator class.) The default behavior of the above operators is the same as for IS [ NOT ] DISTINCT FROM for row constructors (see Section 9.23.5).

To support matching of rows which include elements without a default B-tree operator class, the following operators are defined for composite type comparison: *=, *<>, *<, *<=, *>, and *>=. These operators compare the internal binary representation of the two rows. Two rows might have a different binary representation even though comparisons of the two rows with the equality operator is true. The ordering of rows under these comparison operators is deterministic but not otherwise meaningful. These operators are used internally for materialized views and might be useful for other specialized purposes such as replication but are not intended to be generally useful for writing queries.

### 9.24. Set Returning Functions

This section describes functions that possibly return more than one row. The most widely used functions in this class are series generating functions, as detailed in Table 9.58 and Table 9.59. Other, more specialized set-returning functions are described elsewhere in this manual. See Section 7.2.1.4 for ways to combine multiple set-returning functions.

#### Table 9.58. Series Generating Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>generate_series(...)</td>
<td>int, bigint or numeric</td>
<td>setof int, setof bigint, set of numeric (same as argument type) from start to stop with a step size of one</td>
<td></td>
</tr>
<tr>
<td>generate_series(...)</td>
<td>int, bigint or numeric</td>
<td>setof int, setof bigint or set of numeric (same as argument type) from start to stop with a step size of step</td>
<td></td>
</tr>
<tr>
<td>generate_series(...)</td>
<td>timestamp or timestamp with timezone</td>
<td>setof timestamp or set of timestamp with time zone (same as argument type) from start to stop with a step size of step</td>
<td></td>
</tr>
</tbody>
</table>

When step is positive, zero rows are returned if start is greater than stop. Conversely, when step is negative, zero rows are returned if start is less than stop. Zero rows are also returned for NULL inputs. It is an error for step to be zero. Some examples follow:

```sql
SELECT * FROM generate_series(2,4);
```

```
<table>
<thead>
<tr>
<th>generate_series</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
</tr>
<tr>
<td>2</td>
</tr>
<tr>
<td>3</td>
</tr>
<tr>
<td>4</td>
</tr>
</tbody>
</table>
```

(3 rows)
SELECT * FROM generate_series(5,1,-2);
generate_series
-----------------
   5
   3
   1
(3 rows)

SELECT * FROM generate_series(4,3);
generate_series
-----------------
(0 rows)

SELECT generate_series(1.1, 4, 1.3);
generate_series
-----------------
   1.1
   2.4
   3.7
(3 rows)

-- this example relies on the date-plus-integer operator
SELECT current_date + s.a AS dates FROM generate_series(0,14,7) AS 
s(a); 
dates
-----------
2004-02-05
2004-02-12
2004-02-19
(3 rows)

SELECT * FROM generate_series('2008-03-01 00:00'::timestamp, '2008-03-04 12:00', '10 hours');
generate_series
---------------------
2008-03-01 00:00:00
2008-03-01 10:00:00
2008-03-01 20:00:00
2008-03-02 06:00:00
2008-03-02 16:00:00
2008-03-03 02:00:00
2008-03-03 12:00:00
2008-03-03 22:00:00
2008-03-04 08:00:00
(9 rows)

Table 9.59. Subscript Generating Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>generate_subscript-s</td>
<td>setof int</td>
<td>Generate a series comprising the given array's subscripts.</td>
</tr>
<tr>
<td>(array anyarray, dim int)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>generate_subscript-s</td>
<td>setof int</td>
<td>Generate a series comprising the given array's subscripts. When reverse is true, the series is returned in reverse order.</td>
</tr>
<tr>
<td>(array anyarray, dim int,</td>
<td></td>
<td></td>
</tr>
<tr>
<td>reverse boolean)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
generate_subscripts is a convenience function that generates the set of valid subscripts for the specified dimension of the given array. Zero rows are returned for arrays that do not have the requested dimension, or for NULL arrays (but valid subscripts are returned for NULL array elements). Some examples follow:

```sql
-- basic usage
SELECT generate_subscripts('{NULL,1,NULL,2}'::int[], 1) AS s;

s
---
1
2
3
4
(4 rows)

-- presenting an array, the subscript and the subscripted value requires a subquery
SELECT * FROM arrays;

a
-------------------
{-1,-2}  
{100,200,300}  
(2 rows)

SELECT a AS array, s AS subscript, a[s] AS value
FROM (SELECT generate_subscripts(a, 1) AS s, a FROM arrays) foo;

array   | subscript | value
---------------+-----------+-------
{-1,-2}       |         1 |    -1
{-1,-2}       |         2 |    -2
{100,200,300} |         1 |   100
{100,200,300} |         2 |   200
{100,200,300} |         3 |   300
(5 rows)

-- unnest a 2D array
CREATE OR REPLACE FUNCTION unnest2(anyarray)
RETURNS SETOF anyelement AS $$
select $1[i][j]
from generate_subscripts($1,1) g1(i),
generate_subscripts($1,2) g2(j);
$$ LANGUAGE sql IMMUTABLE;
CREATE FUNCTION
SELECT * FROM unnest2(ARRAY[[1,2],[3,4]]);
unnest2
---------
1
2
3
4
(4 rows)
```

When a function in the FROM clause is suffixed by WITH ORDINALITY, a bigint column is appended to the output which starts from 1 and increments by 1 for each row of the function's output. This is most useful in the case of set returning functions such as unnest().
9.25. System Information Functions

Table 9.60 shows several functions that extract session and system information.

In addition to the functions listed in this section, there are a number of functions related to the statistics system that also provide system information. See Section 28.2.2 for more information.

Table 9.60. Session Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>current_catalog</td>
<td>name</td>
<td>name of current database (called “catalog” in the SQL standard)</td>
</tr>
<tr>
<td>current_database()</td>
<td>name</td>
<td>name of current database</td>
</tr>
<tr>
<td>current_query()</td>
<td>text</td>
<td>text of the currently executing query, as submitted by the client (might contain more than one statement)</td>
</tr>
<tr>
<td>current_role</td>
<td>name</td>
<td>equivalent to current_user</td>
</tr>
<tr>
<td>current_schema()</td>
<td>name</td>
<td>name of current schema</td>
</tr>
<tr>
<td>current_schemas( boolean )</td>
<td>name[]</td>
<td>names of schemas in search path, optionally including implicit schemas</td>
</tr>
<tr>
<td>current_user</td>
<td>name</td>
<td>user name of current execution context</td>
</tr>
<tr>
<td>inet_client_addr()</td>
<td>inet</td>
<td>address of the remote connection</td>
</tr>
<tr>
<td>inet_client_port()</td>
<td>int</td>
<td>port of the remote connection</td>
</tr>
<tr>
<td>inet_server_addr()</td>
<td>inet</td>
<td>address of the local connection</td>
</tr>
<tr>
<td>inet_server_port()</td>
<td>int</td>
<td>port of the local connection</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------------</td>
<td>------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_backend_pid()</code></td>
<td>int</td>
<td>Process ID of the server process attached to the current session</td>
</tr>
<tr>
<td><code>pg_blocking_pids(int)</code></td>
<td>int[]</td>
<td>Process ID(s) that are blocking specified server process ID from acquiring a lock</td>
</tr>
<tr>
<td><code>pg_conf_load_time()</code></td>
<td>timestamp with time zone</td>
<td>configuration load time</td>
</tr>
<tr>
<td><code>pg_current_logfile([text])</code></td>
<td>text</td>
<td>Primary log file name, or log in the requested format, currently in use by the logging collector</td>
</tr>
<tr>
<td><code>pg_my_temp_schema()</code></td>
<td>oid</td>
<td>OID of session's temporary schema, or 0 if none</td>
</tr>
<tr>
<td><code>pg_is_other_temp_schema(oid)</code></td>
<td>boolean</td>
<td>is schema another session's temporary schema?</td>
</tr>
<tr>
<td><code>pg_jit_available()</code></td>
<td>boolean</td>
<td>is JIT compilation available in this session (see Chapter 32)? Returns false if jit is set to false.</td>
</tr>
<tr>
<td><code>pg_listening_channels()</code></td>
<td>setof text</td>
<td>channel names that the session is currently listening on</td>
</tr>
<tr>
<td><code>pg_notification_queue_usage()</code></td>
<td>double</td>
<td>fraction of the asynchronous notification queue currently occupied (0-1)</td>
</tr>
<tr>
<td><code>pg_postmaster_start_time()</code></td>
<td>timestamp with time zone</td>
<td>server start time</td>
</tr>
<tr>
<td><code>pg_safe_snapshot_blocking_pids(int)</code></td>
<td>int[]</td>
<td>Process ID(s) that are blocking specified server process ID from acquiring a safe snapshot</td>
</tr>
<tr>
<td><code>pg_trigger_depth()</code></td>
<td>int</td>
<td>current nesting level of PostgreSQL triggers (0 if not called, directly or indirectly, from inside a trigger)</td>
</tr>
<tr>
<td><code>session_user</code></td>
<td>name</td>
<td>session user name</td>
</tr>
<tr>
<td><code>user</code></td>
<td>name</td>
<td>equivalent to <code>current_user</code></td>
</tr>
<tr>
<td><code>version()</code></td>
<td>text</td>
<td>PostgreSQL version information. See also <code>server_version_num</code> for a machine-readable version.</td>
</tr>
</tbody>
</table>

**Note**

current_catalog, current_role, current_schema, current_user, session_user, and user have special syntactic status in SQL; they must be called without trailing parentheses. (In PostgreSQL, parentheses can optionally be used with current_schema, but not with the others.)

The `session_user` is normally the user who initiated the current database connection; but superusers can change this setting with SET SESSION AUTHORIZATION. The `current_user` is the user identifier that is applicable for permission checking. Normally it is equal to the session user, but
it can be changed with SET ROLE. It also changes during the execution of functions with the attribute
SECURITY DEFINER. In Unix parlance, the session user is the “real user” and the current user
is the “effective user”. current_role and user are synonyms for current_user. (The SQL
standard draws a distinction between current_role and current_user, but PostgreSQL does
not, since it unifies users and roles into a single kind of entity.)

current_schema returns the name of the schema that is first in the search path (or a null value if
the search path is empty). This is the schema that will be used for any tables or other named objects that
are created without specifying a target schema. current_schemas(boolean) returns an array
of the names of all schemas presently in the search path. The Boolean option determines whether or not
implicitly included system schemas such as pg_catalog are included in the returned search path.

The search path can be altered at run time. The command is:

```
SET search_path TO schema [, schema, ...]
```

inet_client_addr returns the IP address of the current client, and inet_client_port re-
turns the port number. inet_server_addr returns the IP address on which the server accepted the
current connection, and inet_server_port returns the port number. All these functions return
NULL if the current connection is via a Unix-domain socket.

pg_blocking_pids returns an array of the process IDs of the sessions that are blocking the server
process with the specified process ID, or an empty array if there is no such server process or it is not
blocked. One server process blocks another if it either holds a lock that conflicts with the blocked
process’s lock request (hard block), or is waiting for a lock that would conflict with the blocked
process’s lock request and is ahead of it in the wait queue (soft block). When using parallel queries the
result always lists client-visible process IDs (that is, pg_backend_pid results) even if the actual
lock is held or awaited by a child worker process. As a result of that, there may be duplicated PIDs in
the result. Also note that when a prepared transaction holds a conflicting lock, it will be represented
by a zero process ID in the result of this function. Frequent calls to this function could have some
impact on database performance, because it needs exclusive access to the lock manager’s shared state
for a short time.

pg_conf_load_time returns the timestamp with time zone when the server configura-
tion files were last loaded. (If the current session was alive at the time, this will be the time when
the session itself re-read the configuration files, so the reading will vary a little in different sessions.
Otherwise it is the time when the postmaster process re-read the configuration files.)

pg_current_logfile returns, as text, the path of the log file(s) currently in use by the logging
collector. The path includes the log_directory directory and the log file name. Log collection must
be enabled or the return value is NULL. When multiple log files exist, each in a different format,
pg_current_logfile called without arguments returns the path of the file having the first format
found in the ordered list: stderr, csvlog. NULL is returned when no log file has any of these formats.
To request a specific file format supply, as text, either csvlog or stderr as the value of the optional
parameter. The return value is NULL when the log format requested is not a configured log_destination.
The pg_current_logfiles reflects the contents of the current_logfiles file.

pg_my_temp_schema returns the OID of the current session’s temporary schema, or zero if it has
none (because it has not created any temporary tables). pg_is_other_temp_schema returns true
if the given OID is the OID of another session’s temporary schema. (This can be useful, for example,
to exclude other sessions’ temporary tables from a catalog display.)
pg_listening_channels returns a set of names of asynchronous notification channels that the current session is listening to. pg_notification_queue_usage returns the fraction of the total available space for notifications currently occupied by notifications that are waiting to be processed, as a double in the range 0-1. See LISTEN and NOTIFY for more information.

pg_postmaster_start_time returns the timestamp with time zone when the server started.

pg_safe_snapshot_blocking_pids returns an array of the process IDs of the sessions that are blocking the server process with the specified process ID from acquiring a safe snapshot, or an empty array if there is no such server process or it is not blocked. A session running a SERIALIZABLE transaction blocks a SERIALizable READ ONLY DEFERRABLE transaction from acquiring a snapshot until the latter determines that it is safe to avoid taking any predicate locks. See Section 13.2.3 for more information about serializable and deferrable transactions. Frequent calls to this function could have some impact on database performance, because it needs access to the predicate lock manager's shared state for a short time.

version returns a string describing the PostgreSQL server's version. You can also get this information from server_version or for a machine-readable version, server_version_num. Software developers should use server_version_num (available since 8.2) or PQserverVersion instead of parsing the text version.

Table 9.61 lists functions that allow the user to query object access privileges programmatically. See Section 5.6 for more information about privileges.

### Table 9.61. Access Privilege Inquiry Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>has_any_column_privilege(user, table, privilege)</td>
<td>boolean</td>
<td>does user have privilege for any column of table</td>
</tr>
<tr>
<td>has_any_column_privilege(table, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for any column of table</td>
</tr>
<tr>
<td>has_column_privilege(user, table, column, privilege)</td>
<td>boolean</td>
<td>does user have privilege for column</td>
</tr>
<tr>
<td>has_column_privilege(table, column, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for column</td>
</tr>
<tr>
<td>has_database_privilege(user, database, privilege)</td>
<td>boolean</td>
<td>does user have privilege for database</td>
</tr>
<tr>
<td>has_database_privilege(database, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for database</td>
</tr>
<tr>
<td>has_foreign_data_wrapper_privilege(user, fdw, privilege)</td>
<td>boolean</td>
<td>does user have privilege for foreign-data wrapper</td>
</tr>
<tr>
<td>has_foreign_data_wrapper_privilege(fdw, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for foreign-data wrapper</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------------------</td>
<td>-------------</td>
<td>--------------------------------------------------</td>
</tr>
<tr>
<td>has_function_privilege(user, function, privilege)</td>
<td>boolean</td>
<td>does user have privilege for function</td>
</tr>
<tr>
<td>has_function_privilege(function, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for function</td>
</tr>
<tr>
<td>has_language_privilege(user, language, privilege)</td>
<td>boolean</td>
<td>does user have privilege for language</td>
</tr>
<tr>
<td>has_language_privilege(language, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for language</td>
</tr>
<tr>
<td>has_schema_privilege(user, schema, privilege)</td>
<td>boolean</td>
<td>does user have privilege for schema</td>
</tr>
<tr>
<td>has_schema_privilege(schema, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for schema</td>
</tr>
<tr>
<td>has_sequence_privilege(user, sequence, privilege)</td>
<td>boolean</td>
<td>does user have privilege for sequence</td>
</tr>
<tr>
<td>has_sequence_privilege(sequence, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for sequence</td>
</tr>
<tr>
<td>has_server_privilege(user, server, privilege)</td>
<td>boolean</td>
<td>does user have privilege for foreign server</td>
</tr>
<tr>
<td>has_server_privilege(server, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for foreign server</td>
</tr>
<tr>
<td>has_table_privilege(user, table, privilege)</td>
<td>boolean</td>
<td>does user have privilege for table</td>
</tr>
<tr>
<td>has_table_privilege(table, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for table</td>
</tr>
<tr>
<td>has_tablespace_privilege(user, tablespace, privilege)</td>
<td>boolean</td>
<td>does user have privilege for tablespace</td>
</tr>
<tr>
<td>has_tablespace_privilege(tablespace, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for tablespace</td>
</tr>
<tr>
<td>has_type_privilege(user, type, privilege)</td>
<td>boolean</td>
<td>does user have privilege for type</td>
</tr>
<tr>
<td>has_type_privilege(type, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for type</td>
</tr>
<tr>
<td>pg_has_role(user, role, privilege)</td>
<td>boolean</td>
<td>does user have privilege for role</td>
</tr>
</tbody>
</table>
has_table_privilege checks whether a user can access a table in a particular way. The user can be specified by name, by OID (pg_authid.oid), public to indicate the PUBLIC pseudo-role, or if the argument is omitted current_user is assumed. The table can be specified by name or by OID. (Thus, there are actually six variants of has_table_privilege, which can be distinguished by the number and types of their arguments.) When specifying by name, the name can be schema-qualified if necessary. The desired access privilege type is specified by a text string, which must evaluate to one of the values SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, or TRIGGER. Optionally, WITH GRANT OPTION can be added to a privilege type to test whether the privilege is held with grant option. Also, multiple privilege types can be listed separated by commas, in which case the result will be true if any of the listed privileges is held. (Case of the privilege string is not significant, and extra whitespace is allowed between but not within privilege names.) Some examples:

```sql
SELECT has_table_privilege('myschema.mytable', 'select');
SELECT has_table_privilege('joe', 'mytable', 'INSERT, SELECT WITH
GRANT OPTION');
```

has_sequence_privilege checks whether a user can access a sequence in a particular way. The possibilities for its arguments are analogous to has_table_privilege. The desired access privilege type must evaluate to one of USAGE, SELECT, or UPDATE.

has_any_column_privilege checks whether a user can access any column of a table in a particular way. Its argument possibilities are analogous to has_table_privilege, except that the desired access privilege type must evaluate to some combination of SELECT, INSERT, UPDATE, or REFERENCES. Note that having any of these privileges at the table level implicitly grants it for each column of the table, so has_any_column_privilege will always return true if has_table_privilege does for the same arguments. But has_any_column_privilege also succeeds if there is a column-level grant of the privilege for at least one column.

has_column_privilege checks whether a user can access a column in a particular way. Its argument possibilities are analogous to has_table_privilege, with the addition that the column can be specified either by name or attribute number. The desired access privilege type must evaluate to some combination of SELECT, INSERT, UPDATE, or REFERENCES. Note that having any of these privileges at the table level implicitly grants it for each column of the table.

has_database_privilege checks whether a user can access a database in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to some combination of CREATE, CONNECT, TEMPORARY, or TEMP (which is equivalent to TEMPORARY).

has_function_privilege checks whether a user can access a function in a particular way. Its argument possibilities are analogous to has_table_privilege. When specifying a function by a text string rather than by OID, the allowed input is the same as for the regprocedure data type (see Section 8.19). The desired access privilege type must evaluate to EXECUTE. An example is:

```sql
SELECT has_function_privilege('joeuser', 'myfunc(int, text)', 'execute');
```

has_foreign_data_wrapper_privilege checks whether a user can access a foreign-data wrapper in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to USAGE.
has_language_privilege checks whether a user can access a procedural language in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to USAGE.

has_schema_privilege checks whether a user can access a schema in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to some combination of CREATE or USAGE.

has_server_privilege checks whether a user can access a foreign server in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to USAGE.

has_tablespace_privilege checks whether a user can access a tablespace in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to CREATE.

has_type_privilege checks whether a user can access a type in a particular way. Its argument possibilities are analogous to has_table_privilege. When specifying a type by a text string rather than by OID, the allowed input is the same as for the regtype data type (see Section 8.19). The desired access privilege type must evaluate to USAGE.

pg_has_role checks whether a user can access a role in a particular way. Its argument possibilities are analogous to has_table_privilege, except that public is not allowed as a user name. The desired access privilege type must evaluate to some combination of MEMBER or USAGE. MEMBER denotes direct or indirect membership in the role (that is, the right to do SET ROLE), while USAGE denotes whether the privileges of the role are immediately available without doing SET ROLE.

row_security_active checks whether row level security is active for the specified table in the context of the current_user and environment. The table can be specified by name or by OID.

Table 9.62 shows functions that determine whether a certain object is visible in the current schema search path. For example, a table is said to be visible if its containing schema is in the search path and no table of the same name appears earlier in the search path. This is equivalent to the statement that the table can be referenced by name without explicit schema qualification. To list the names of all visible tables:

SELECT relname FROM pg_class WHERE pg_table_is_visible(oid);

Table 9.62. Schema Visibility Inquiry Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_collation_is_visible(collation_oid)</td>
<td>boolean</td>
<td>is collation visible in search path</td>
</tr>
<tr>
<td>pg_conversion_is_visible(conversion_oid)</td>
<td>boolean</td>
<td>is conversion visible in search path</td>
</tr>
<tr>
<td>pg_function_is_visible(function_oid)</td>
<td>boolean</td>
<td>is function visible in search path</td>
</tr>
<tr>
<td>pg_opclass_is_visible(opclass_oid)</td>
<td>boolean</td>
<td>is operator class visible in search path</td>
</tr>
<tr>
<td>pg_operator_is_visible(operator_oid)</td>
<td>boolean</td>
<td>is operator visible in search path</td>
</tr>
<tr>
<td>pg_opfamily_is_visible(opclass_oid)</td>
<td>boolean</td>
<td>is operator family visible in search path</td>
</tr>
<tr>
<td>pg_statistics_obj_is_visible(stat_hash oid)</td>
<td>boolean</td>
<td>is statistics object visible in search path</td>
</tr>
</tbody>
</table>
Each function performs the visibility check for one type of database object. Note that `pg_table_is_visible` can also be used with views, materialized views, indexes, sequences and foreign tables; `pg_function_is_visible` can also be used with procedures and aggregates; `pg_type_is_visible` can also be used with domains. For functions and operators, an object in the search path is visible if there is no object of the same name and argument data type(s) earlier in the path. For operator classes, both name and associated index access method are considered.

All these functions require object OIDs to identify the object to be checked. If you want to test an object by name, it is convenient to use the OID alias types (`regclass`, `regtype`, `regprocedure`, `regoperator`, `regconfig`, or `regdictionary`), for example:

```
SELECT pg_type_is_visible('myschema.widget'::regtype);
```

Note that it would not make much sense to test a non-schema-qualified type name in this way — if the name can be recognized at all, it must be visible.

Table 9.63 lists functions that extract information from the system catalogs.

### Table 9.63. System Catalog Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>format_type(type_oid, typemod)</code></td>
<td>text</td>
<td>get SQL name of a data type</td>
</tr>
<tr>
<td><code>pg_get_constraintdef(constraint_oid)</code></td>
<td>text</td>
<td>get definition of a constraint</td>
</tr>
<tr>
<td><code>pg_get_constraintdef(constraint_oid, pretty_bool)</code></td>
<td>text</td>
<td>get definition of a constraint</td>
</tr>
<tr>
<td><code>pg_get_expr(pg_node_tree, relation_oid)</code></td>
<td>text</td>
<td>decompile internal form of an expression, assuming that any Vars in it refer to the relation indicated by the second parameter</td>
</tr>
<tr>
<td><code>pg_get_expr(pg_node_tree, relation_oid, pretty_bool)</code></td>
<td>text</td>
<td>decompile internal form of an expression, assuming that any Vars in it refer to the relation indicated by the second parameter</td>
</tr>
<tr>
<td><code>pg_get_functiondef(func_oid)</code></td>
<td>text</td>
<td>get definition of a function or procedure</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>------</td>
<td>-------------</td>
<td>-------------</td>
</tr>
<tr>
<td>pg_get_function_arguments(func_oid)</td>
<td>text</td>
<td>get argument list of function's or procedure's definition (with default values)</td>
</tr>
<tr>
<td>pg_get_function_identity_arguments(func_oid)</td>
<td>text</td>
<td>get argument list to identify a function or procedure (without default values)</td>
</tr>
<tr>
<td>pg_get_function_result(func_oid)</td>
<td>text</td>
<td>get RETURNS clause for function (returns null for a procedure)</td>
</tr>
<tr>
<td>pg_get_indexdef(index_oid)</td>
<td>text</td>
<td>get CREATE_INDEX command for index</td>
</tr>
<tr>
<td>pg_get_indexdef(index_oid, column_no, pretty_bool)</td>
<td>text</td>
<td>get CREATE_INDEX command for index, or definition of just one index column when column_no is not zero</td>
</tr>
<tr>
<td>pg_get_keywords()</td>
<td>setof record</td>
<td>get list of SQL keywords and their categories</td>
</tr>
<tr>
<td>pg_get_ruledef(rule_oid)</td>
<td>text</td>
<td>get CREATE RULE command for rule</td>
</tr>
<tr>
<td>pg_get_ruledef(rule_oid, pretty_bool)</td>
<td>text</td>
<td>get CREATE RULE command for rule</td>
</tr>
<tr>
<td>pg_get_serial_sequence(table_name, column_name)</td>
<td>text</td>
<td>get name of the sequence that a serial or identity column uses</td>
</tr>
<tr>
<td>pg_get_statisticsobject(statobj_oid)</td>
<td>text</td>
<td>get CREATE STATISTICS command for extended statistics object</td>
</tr>
<tr>
<td>pg_get_triggerdef(trigger_oid)</td>
<td>text</td>
<td>get CREATE [ CONSTRAINT ] TRIGGER command for trigger</td>
</tr>
<tr>
<td>pg_get_triggerdef(trigger_oid, pretty_bool)</td>
<td>text</td>
<td>get CREATE [ CONSTRAINT ] TRIGGER command for trigger</td>
</tr>
<tr>
<td>pg_get_userbyid(role_oid)</td>
<td>name</td>
<td>get role name with given OID</td>
</tr>
<tr>
<td>pg_get_viewdef(view_name)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view (deprecated)</td>
</tr>
<tr>
<td>pg_get_viewdef(view_name, pretty_bool)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view (deprecated)</td>
</tr>
<tr>
<td>pg_get_viewdef(view_oid)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view</td>
</tr>
<tr>
<td>pg_get_viewdef(view_oid, pretty_bool)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>------</td>
<td>-------------</td>
<td>-------------</td>
</tr>
<tr>
<td>pg_get_viewdef(view_oid, wrap_column_int)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view; lines with fields are wrapped to specified number of columns, pretty-printing is implied</td>
</tr>
<tr>
<td>pg_index_column_has_property(index_oid, column_no, prop_name)</td>
<td>boolean</td>
<td>test whether an index column has a specified property</td>
</tr>
<tr>
<td>pg_index_has_property(index_oid, prop_name)</td>
<td>boolean</td>
<td>test whether an index has a specified property</td>
</tr>
<tr>
<td>pg_indexam_has_property(am_oid, prop_name)</td>
<td>boolean</td>
<td>test whether an index access method has a specified property</td>
</tr>
<tr>
<td>pg_options_to_table(reloptions)</td>
<td>setof record</td>
<td>get the set of storage option name/value pairs</td>
</tr>
<tr>
<td>pg_tablespace_databases(tablespace_oid)</td>
<td>setof oid</td>
<td>get the set of database OIDs that have objects in the tablespace</td>
</tr>
<tr>
<td>pg_tablespace_location(tablespace_oid)</td>
<td>text</td>
<td>get the path in the file system that this tablespace is located in</td>
</tr>
<tr>
<td>pg_typeof(any)</td>
<td>regtype</td>
<td>get the data type of any value</td>
</tr>
<tr>
<td>collation for (any)</td>
<td>text</td>
<td>get the collation of the argument</td>
</tr>
<tr>
<td>to_regclass(rel_name)</td>
<td>regclass</td>
<td>get the OID of the named relation</td>
</tr>
<tr>
<td>to_regproc(func_name)</td>
<td>regproc</td>
<td>get the OID of the named function</td>
</tr>
<tr>
<td>to_regprocedure(func_name)</td>
<td>regprocedure</td>
<td>get the OID of the named function</td>
</tr>
<tr>
<td>to_regoperator(operator_name)</td>
<td>regoperator</td>
<td>get the OID of the named operator</td>
</tr>
<tr>
<td>to_regoperator(operator_name)</td>
<td>regoperator</td>
<td>get the OID of the named operator</td>
</tr>
<tr>
<td>to_regtype(type_name)</td>
<td>regtype</td>
<td>get the OID of the named type</td>
</tr>
<tr>
<td>to_regnamespace(schema_name)</td>
<td>regnamespace</td>
<td>get the OID of the named schema</td>
</tr>
<tr>
<td>to_regrole(role_name)</td>
<td>regrole</td>
<td>get the OID of the named role</td>
</tr>
</tbody>
</table>

format_type returns the SQL name of a data type that is identified by its type OID and possibly a type modifier. Pass NULL for the type modifier if no specific modifier is known.

pg_get_keywords returns a set of records describing the SQL keywords recognized by the server. The word column contains the keyword. The catcode column contains a category code: U for unreserved, C for column name, T for type or function name, or R for reserved. The catdesc column contains a possibly-localized string describing the category.

pg_get_constraintdef, pg_get_indexdef, pg_get_ruledef, pg_get_statisticsobjdef, and pg_get_triggerdef, respectively reconstruct the creating command for a constraint, index, rule, extended statistics object, or trigger. (Note that this is a decompiled reconstruction, not the original text of the command.) pg_get_expr decompiles the internal form of an
individual expression, such as the default value for a column. It can be useful when examining the contents of system catalogs. If the expression might contain Vars, specify the OID of the relation they refer to as the second parameter; if no Vars are expected, zero is sufficient. `pg_get_viewdef` reconstructs the SELECT query that defines a view. Most of these functions come in two variants, one of which can optionally "pretty-print" the result. The pretty-printed format is more readable, but the default format is more likely to be interpreted the same way by future versions of PostgreSQL; avoid using pretty-printed output for dump purposes. Passing `false` for the pretty-print parameter yields the same result as the variant that does not have the parameter at all.

`pg_get_functiondef` returns a complete CREATE OR REPLACE FUNCTION statement for a function. `pg_get_function_arguments` returns the argument list of a function, in the form it would need to appear in within CREATE FUNCTION. `pg_get_function_result` similarly returns the appropriate RETURNS clause for the function. `pg_get_function_identity_arguments` returns the argument list necessary to identify a function, in the form it would need to appear in within ALTER FUNCTION, for instance. This form omits default values.

`pg_get_serial_sequence` returns the name of the sequence associated with a column, or NULL if no sequence is associated with the column. If the column is an identity column, the associated sequence is the sequence internally created for the identity column. For columns created using one of the serial types (serial, smallserial, bigserial), it is the sequence created for that serial column definition. In the latter case, this association can be modified or removed with ALTER SEQUENCE OWNED BY. (The function probably should have been called `pg_get_owned_sequence`; its current name reflects the fact that it has typically been used with serial or bigserial columns.) The first input parameter is a table name with optional schema, and the second parameter is a column name. Because the first parameter is potentially a schema and table, it is not treated as a double-quoted identifier, meaning it is lower cased by default, while the second parameter, being just a column name, is treated as double-quoted and has its case preserved. The function returns a value suitably formatted for passing to sequence functions (see Section 9.16). A typical use is in reading the current value of a sequence for an identity or serial column, for example:

```sql
SELECT currval(pg_get_serial_sequence('sometable', 'id'));
```

`pg_get_userbyid` extracts a role's name given its OID.

`pg_index_column_has_property`, `pg_index_has_property`, and `pg_index-am_has_property` return whether the specified index column, index, or index access method possesses the named property. NULL is returned if the property name is not known or does not apply to the particular object, or if the OID or column number does not identify a valid object. Refer to Table 9.64 for column properties, Table 9.65 for index properties, and Table 9.66 for access method properties. (Note that extension access methods can define additional property names for their indexes.)

### Table 9.64. Index Column Properties

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>asc</td>
<td>Does the column sort in ascending order on a forward scan?</td>
</tr>
<tr>
<td>desc</td>
<td>Does the column sort in descending order on a forward scan?</td>
</tr>
<tr>
<td>nulls_first</td>
<td>Does the column sort with nulls first on a forward scan?</td>
</tr>
<tr>
<td>nulls_last</td>
<td>Does the column sort with nulls last on a forward scan?</td>
</tr>
<tr>
<td>orderable</td>
<td>Does the column possess any defined sort ordering?</td>
</tr>
<tr>
<td>Name</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>distance_orderable</td>
<td>Can the column be scanned in order by a “distance” operator, for example ORDER BY col &lt;-&gt; constant?</td>
</tr>
<tr>
<td>returnable</td>
<td>Can the column value be returned by an index-only scan?</td>
</tr>
<tr>
<td>search_array</td>
<td>Does the column natively support col = ANY(array) searches?</td>
</tr>
<tr>
<td>search_nulls</td>
<td>Does the column support IS NULL and IS NOT NULL searches?</td>
</tr>
</tbody>
</table>

**Table 9.65. Index Properties**

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>clusterable</td>
<td>Can the index be used in a CLUSTER command?</td>
</tr>
<tr>
<td>index_scan</td>
<td>Does the index support plain (non-bitmap) scans?</td>
</tr>
<tr>
<td>bitmap_scan</td>
<td>Does the index support bitmap scans?</td>
</tr>
<tr>
<td>backward_scan</td>
<td>Can the scan direction be changed in mid-scan (to support FETCH BACKWARD on a cursor without needing materialization)?</td>
</tr>
</tbody>
</table>

**Table 9.66. Index Access Method Properties**

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>can_order</td>
<td>Does the access method support ASC, DESC and related keywords in CREATE_INDEX?</td>
</tr>
<tr>
<td>can_unique</td>
<td>Does the access method support unique indexes?</td>
</tr>
<tr>
<td>can_multi_col</td>
<td>Does the access method support indexes with multiple columns?</td>
</tr>
<tr>
<td>can_exclude</td>
<td>Does the access method support exclusion constraints?</td>
</tr>
<tr>
<td>can_include</td>
<td>Does the access method support the INCLUDE clause of CREATE_INDEX?</td>
</tr>
</tbody>
</table>

`pg_options_to_table` returns the set of storage option name/value pairs (option_name/option_value) when passed `pg_class.reloptions` or `pg_attribute.attoptions`.

`pg_tablespace_databases` allows a tablespace to be examined. It returns the set of OIDs of databases that have objects stored in the tablespace. If this function returns any rows, the tablespace is not empty and cannot be dropped. To display the specific objects populating the tablespace, you will need to connect to the databases identified by `pg_tablespace_databases` and query their `pg_class` catalogs.

`pg_typeof` returns the OID of the data type of the value that is passed to it. This can be helpful for troubleshooting or dynamically constructing SQL queries. The function is declared as returning `regtype`, which is an OID alias type (see Section 8.19); this means that it is the same as an OID for comparison purposes but displays as a type name. For example:

```sql
SELECT pg_typeof(33);
```

`pg_typeof ___typeof___

332
The expression `collation for` returns the collation of the value that is passed to it. Example:

```sql
SELECT collation for (description) FROM pg_description LIMIT 1;
```

```
pg_collation_for
------------------
"default"
(1 row)
```

SELECT collation for ('foo' COLLATE "de_DE");

```
pg_collation_for
------------------
"de_DE"
(1 row)
```

The value might be quoted and schema-qualified. If no collation is derived for the argument expression, then a null value is returned. If the argument is not of a collatable data type, then an error is raised.

The `to_regclass`, `to_regproc`, `to_regprocedure`, `to_regoperator`, `to_regtype`, `to_regnamespace`, and `to_regrole` functions translate relation, function, operator, type, schema, and role names (given as text) to objects of type `regclass`, `regproc`, `regprocedure`, `regoper`, `regoperator`, `regtype`, `regnamespace`, and `regrole` respectively. These functions differ from a cast from text in that they don’t accept a numeric OID, and that they return null rather than throwing an error if the name is not found (or, for `to_regproc` and `to_regoperator`, if the given name matches multiple objects).

Table 9.67 lists functions related to database object identification and addressing.

### Table 9.67. Object Information and Addressing Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_describe_object(classid oid, objid oid, objsubid integer)</code></td>
<td>text</td>
<td>get description of a database object</td>
</tr>
<tr>
<td><code>pg_identify_object(classid oid, objid oid, objsubid integer)</code></td>
<td><code>type text, schema text, name text, identity text</code></td>
<td>get identity of a database object</td>
</tr>
<tr>
<td><code>pg_identify_object_as_address(classid oid, objid oid, objsubid integer)</code></td>
<td><code>type text, object_names text[], object_args text[]</code></td>
<td>get external representation of a database object’s address</td>
</tr>
<tr>
<td><code>pg_get_object_address(type text, object_names text[], object_args text[])</code></td>
<td><code>classid oid, objid oid, objsubid integer</code></td>
<td>get address of a database object from its external representation</td>
</tr>
</tbody>
</table>

`pg_describe_object` returns a textual description of a database object specified by catalog OID, object OID, and sub-object ID (such as a column number within a table; the sub-object ID is zero...
when referring to a whole object). This description is intended to be human-readable, and might be translated, depending on server configuration. This is useful to determine the identity of an object as stored in the `pg_depend` catalog.

`pg_identify_object` returns a row containing enough information to uniquely identify the database object specified by catalog OID, object OID and sub-object ID. This information is intended to be machine-readable, and is never translated. `type` identifies the type of database object; `schema` is the schema name that the object belongs in, or NULL for object types that do not belong to schemas; `name` is the name of the object, quoted if necessary, if the name (along with schema name, if pertinent) is sufficient to uniquely identify the object, otherwise NULL; `identity` is the complete object identity, with the precise format depending on object type, and each name within the format being schema-qualified and quoted as necessary.

`pg_identify_object_as_address` returns a row containing enough information to uniquely identify the database object specified by catalog OID, object OID and sub-object ID. The returned information is independent of the current server, that is, it could be used to identify an identically named object in another server. `type` identifies the type of database object; `object_names` and `object_args` are text arrays that together form a reference to the object. These three values can be passed to `pg_get_object_address` to obtain the internal address of the object. This function is the inverse of `pg_identify_object_as_address`.

`pg_get_object_address` returns a row containing enough information to uniquely identify the database object specified by its type and object name and argument arrays. The returned values are the ones that would be used in system catalogs such as `pg_depend` and can be passed to other system functions such as `pg_identify_object` or `pg_describe_object`. `classid` is the OID of the system catalog containing the object; `objid` is the OID of the object itself, and `objsubid` is the sub-object ID, or zero if none. This function is the inverse of `pg_identify_object_as_address`.

The functions shown in Table 9.68 extract comments previously stored with the `COMMENT` command. A null value is returned if no comment could be found for the specified parameters.

**Table 9.68. Comment Information Functions**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>col_description(table_oid, column_num)</code></td>
<td>text</td>
<td>get comment for a table column</td>
</tr>
<tr>
<td><code>obj_description(object_oid, catalog_name)</code></td>
<td>text</td>
<td>get comment for a database object</td>
</tr>
<tr>
<td><code>obj_description(object_oid)</code></td>
<td>text</td>
<td>get comment for a database object (deprecated)</td>
</tr>
<tr>
<td><code>shobj_description(object_oid, catalog_name)</code></td>
<td>text</td>
<td>get comment for a shared database object</td>
</tr>
</tbody>
</table>

`col_description` returns the comment for a table column, which is specified by the OID of its table and its column number. (`obj_description` cannot be used for table columns since columns do not have OIDs of their own.)

The two-parameter form of `obj_description` returns the comment for a database object specified by its OID and the name of the containing system catalog. For example, `obj_description(123456, 'pg_class')` would retrieve the comment for the table with OID 123456. The one-parameter form of `obj_description` requires only the object OID. It is deprecated since there is no guarantee that OIDs are unique across different system catalogs; therefore, the wrong comment might be returned.
shobj_description is used just like obj_description except it is used for retrieving comments on shared objects. Some system catalogs are global to all databases within each cluster, and the descriptions for objects in them are stored globally as well.

The functions shown in Table 9.69 provide server transaction information in an exportable form. The main use of these functions is to determine which transactions were committed between two snapshots.

### Table 9.69. Transaction IDs and Snapshots

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>txid_current()</td>
<td>bigint</td>
<td>get current transaction ID, assigning a new one if the current transaction does not have one</td>
</tr>
<tr>
<td>txid_current_if_assigned()</td>
<td>bigint</td>
<td>same as txid_current() but returns null instead of assigning a new transaction ID if none is already assigned</td>
</tr>
<tr>
<td>txid_current_snapshot()</td>
<td>txid_snapshot</td>
<td>get current snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xip(txid_snapshot)</td>
<td>setof bigint</td>
<td>get in-progress transaction IDs in snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xmax(txid_snapshot)</td>
<td>bigint</td>
<td>get xmax of snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xmin(txid_snapshot)</td>
<td>bigint</td>
<td>get xmin of snapshot</td>
</tr>
<tr>
<td>txid_visible_in_snapshot(bigint, txid_snapshot)</td>
<td>boolean</td>
<td>is transaction ID visible in snapshot? (do not use with subtransaction ids)</td>
</tr>
<tr>
<td>txid_status(bigint)</td>
<td>text</td>
<td>report the status of the given transaction: committed, aborted, in progress, or null if the transaction ID is too old</td>
</tr>
</tbody>
</table>

The internal transaction ID type (xid) is 32 bits wide and wraps around every 4 billion transactions. However, these functions export a 64-bit format that is extended with an “epoch” counter so it will not wrap around during the life of an installation. The data type used by these functions, txid_snapshot, stores information about transaction ID visibility at a particular moment in time. Its components are described in Table 9.70.

### Table 9.70. Snapshot Components

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xmin</td>
<td>Earliest transaction ID (xid) that is still active. All earlier transactions will either be committed and visible, or rolled back and dead.</td>
</tr>
<tr>
<td>xmax</td>
<td>First as-yet-unassigned xid. All txids greater than or equal to this are not yet started as of the time of the snapshot, and thus invisible.</td>
</tr>
<tr>
<td>xip_list</td>
<td>Active txids at the time of the snapshot. The list includes only those active txids between xmin and xmax; there might be active txids higher than</td>
</tr>
</tbody>
</table>
Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>xmax</strong></td>
<td>A txid that is xmin &lt;= txid &lt; xmax and not in this list was already completed at the time of the snapshot, and thus either visible or dead according to its commit status. The list does not include txids of subtransactions.</td>
</tr>
</tbody>
</table>

**txid_snapshot**'s textual representation is xmin:xmax:xip_list. For example 10:20:10,14,15 means xmin=10, xmax=20, xip_list=10, 14, 15.

**txid_status**(bigint) reports the commit status of a recent transaction. Applications may use it to determine whether a transaction committed or aborted when the application and database server become disconnected while a COMMIT is in progress. The status of a transaction will be reported as either in progress, committed, or aborted, provided that the transaction is recent enough that the system retains the commit status of that transaction. If is old enough that no references to that transaction survive in the system and the commit status information has been discarded, this function will return NULL. Note that prepared transactions are reported as in progress; applications must check pg_prepared_xacts if they need to determine whether the txid is a prepared transaction.

The functions shown in Table 9.71 provide information about transactions that have been already committed. These functions mainly provide information about when the transactions were committed. They only provide useful data when track_commit_timestamp configuration option is enabled and only for transactions that were committed after it was enabled.

**Table 9.71. Committed transaction information**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_xact_commit_timestamp(xid)</td>
<td>timestamp with time zone</td>
<td>get commit timestamp of a transaction</td>
</tr>
<tr>
<td>pg_last_committed_xact()</td>
<td>xid, xid, timestamp time-stamp with time zone</td>
<td>get transaction ID and commit timestamp of latest committed transaction</td>
</tr>
</tbody>
</table>

The functions shown in Table 9.72 print information initialized during initdb, such as the catalog version. They also show information about write-ahead logging and checkpoint processing. This information is cluster-wide, and not specific to any one database. They provide most of the same information, from the same source, as pg_controldata, although in a form better suited to SQL functions.

**Table 9.72. Control Data Functions**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_control_checkpoint()</td>
<td>record</td>
<td>Returns information about current checkpoint state.</td>
</tr>
<tr>
<td>pg_control_system()</td>
<td>record</td>
<td>Returns information about current control file state.</td>
</tr>
<tr>
<td>pg_control_init()</td>
<td>record</td>
<td>Returns information about cluster initialization state.</td>
</tr>
<tr>
<td>pg_control_recovery()</td>
<td>record</td>
<td>Returns information about recovery state.</td>
</tr>
</tbody>
</table>

pg_control_checkpoint returns a record, shown in Table 9.73

**Table 9.73. pg_control_checkpoint Columns**

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>checkpoint_lsn</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>Column Name</td>
<td>Data Type</td>
</tr>
<tr>
<td>---------------------</td>
<td>----------------------</td>
</tr>
<tr>
<td>redo_lsn</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>redo_wal_file</td>
<td>text</td>
</tr>
<tr>
<td>timeline_id</td>
<td>integer</td>
</tr>
<tr>
<td>prev_timeline_id</td>
<td>integer</td>
</tr>
<tr>
<td>full_page_writes</td>
<td>boolean</td>
</tr>
<tr>
<td>next_xid</td>
<td>text</td>
</tr>
<tr>
<td>next_oid</td>
<td>oid</td>
</tr>
<tr>
<td>next_multixact_id</td>
<td>xid</td>
</tr>
<tr>
<td>next_multi_offset</td>
<td>xid</td>
</tr>
<tr>
<td>oldest_xid</td>
<td>xid</td>
</tr>
<tr>
<td>oldest_xid_dbid</td>
<td>oid</td>
</tr>
<tr>
<td>oldest_active_xid</td>
<td>xid</td>
</tr>
<tr>
<td>oldest_multi_xid</td>
<td>xid</td>
</tr>
<tr>
<td>oldest_multi_dbid</td>
<td>oid</td>
</tr>
<tr>
<td>oldest_commit_ts_xid</td>
<td>xid</td>
</tr>
<tr>
<td>newest_commit_ts_xid</td>
<td>xid</td>
</tr>
<tr>
<td>checkpoint_time</td>
<td>timestamp with time zone</td>
</tr>
</tbody>
</table>

`pg_control_system` returns a record, shown in Table 9.74

**Table 9.74. pg_control_system Columns**

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_control_version</td>
<td>integer</td>
</tr>
<tr>
<td>catalog_version_no</td>
<td>integer</td>
</tr>
<tr>
<td>system_identifier</td>
<td>bigint</td>
</tr>
<tr>
<td>pg_control_last_modified</td>
<td>timestamp with time zone</td>
</tr>
</tbody>
</table>

`pg_control_init` returns a record, shown in Table 9.75

**Table 9.75. pg_control_init Columns**

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>max_data_alignment</td>
<td>integer</td>
</tr>
<tr>
<td>database_block_size</td>
<td>integer</td>
</tr>
<tr>
<td>blocks_per_segment</td>
<td>integer</td>
</tr>
<tr>
<td>wal_block_size</td>
<td>integer</td>
</tr>
<tr>
<td>bytes_per_wal_segment</td>
<td>integer</td>
</tr>
<tr>
<td>max_identifier_length</td>
<td>integer</td>
</tr>
<tr>
<td>max_index_columns</td>
<td>integer</td>
</tr>
<tr>
<td>max_toast_chunk_size</td>
<td>integer</td>
</tr>
<tr>
<td>large_object_chunk_size</td>
<td>integer</td>
</tr>
<tr>
<td>float4_pass_by_value</td>
<td>boolean</td>
</tr>
<tr>
<td>float8_pass_by_value</td>
<td>boolean</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_page_checksum_version</td>
<td>integer</td>
</tr>
</tbody>
</table>

`pg_control_recovery` returns a record, shown in Table 9.76

#### Table 9.76. `pg_control_recovery` Columns

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>min_recovery_end_lsn</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>min_recovery_end_timeline</td>
<td>integer</td>
</tr>
<tr>
<td>backup_start_lsn</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>backup_end_lsn</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>end_of_backup_record_required</td>
<td>boolean</td>
</tr>
</tbody>
</table>

### 9.26. System Administration Functions

The functions described in this section are used to control and monitor a PostgreSQL installation.

#### 9.26.1. Configuration Settings Functions

Table 9.77 shows the functions available to query and alter run-time configuration parameters.

#### Table 9.77. Configuration Settings Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>current_setting(setting_name [, missing_ok ] )</td>
<td>text</td>
<td>get current value of setting</td>
</tr>
<tr>
<td>set_config(setting_name, new_value, is_local)</td>
<td>text</td>
<td>set parameter and return new value</td>
</tr>
</tbody>
</table>

The function `current_setting` yields the current value of the setting `setting_name`. It corresponds to the SQL command `SHOW`. An example:

```
SELECT current_setting('datestyle');
```

```
current_setting
--------------
ISO, MDY
(1 row)
```

If there is no setting named `setting_name`, `current_setting` throws an error unless `missing_ok` is supplied and is true.

`set_config` sets the parameter `setting_name` to `new_value`. If `is_local` is `true`, the new value will only apply to the current transaction. If you want the new value to apply for the current session, use `false` instead. The function corresponds to the SQL command `SET`. An example:

```
SELECT set_config('log_statement_stats', 'off', false);
```
9.26.2. Server Signaling Functions

The functions shown in Table 9.78 send control signals to other server processes. Use of these functions is restricted to superusers by default but access may be granted to others using GRANT, with noted exceptions.

Table 9.78. Server Signaling Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_cancel_backend(pid int)</td>
<td>boolean</td>
<td>Cancel a backend's current query. This is also allowed if the calling role is a member of the role whose backend is being canceled or the calling role has been granted pg_signal_backend, however only superusers can cancel superuser backends.</td>
</tr>
<tr>
<td>pg_reload_conf()</td>
<td>boolean</td>
<td>Cause server processes to reload their configuration files</td>
</tr>
<tr>
<td>pg_rotate_logfile()</td>
<td>boolean</td>
<td>Rotate server's log file</td>
</tr>
<tr>
<td>pg_terminate_backend(pid int)</td>
<td>boolean</td>
<td>Terminate a backend. This is also allowed if the calling role is a member of the role whose backend is being terminated or the calling role has been granted pg_signal_backend, however only superusers can terminate superuser backends.</td>
</tr>
</tbody>
</table>

Each of these functions returns true if successful and false otherwise.

pg_cancel_backend and pg_terminate_backend send signals (SIGINT or SIGTERM respectively) to backend processes identified by process ID. The process ID of an active backend can be found from the pid column of the pg_stat_activity view, or by listing the postgres processes on the server (using ps on Unix or the Task Manager on Windows). The role of an active backend can be found from the usename column of the pg_stat_activity view.

pg_reload_conf sends a SIGHUP signal to the server, causing configuration files to be reloaded by all server processes.

pg_rotate_logfile signals the log-file manager to switch to a new output file immediately. This works only when the built-in log collector is running, since otherwise there is no log-file manager subprocess.

9.26.3. Backup Control Functions

The functions shown in Table 9.79 assist in making on-line backups. These functions cannot be executed during recovery (except pg_is_in_backup, pg_backup_start_time and pg_wal_lsn_diff).
Table 9.79. Backup Control Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_create_restore_point(name text)</td>
<td>pg_lsn</td>
<td>Create a named point for performing restore (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td>pg_current_wal_flush_lsn()</td>
<td>pg_lsn</td>
<td>Get current write-ahead log flush location</td>
</tr>
<tr>
<td>pg_current_wal_insert_lsn()</td>
<td>pg_lsn</td>
<td>Get current write-ahead log insert location</td>
</tr>
<tr>
<td>pg_current_wal_lsn()</td>
<td>pg_lsn</td>
<td>Get current write-ahead log write location</td>
</tr>
<tr>
<td>pg_start_backup(label text [, fast boolean [, exclusive boolean ]])</td>
<td>pg_lsn</td>
<td>Prepare for performing on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td>pg_stop_backup()</td>
<td>pg_lsn</td>
<td>Finish performing exclusive on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td>pg_stop_backup(exclusive boolean [, wait_for_archive boolean ])</td>
<td>setof record</td>
<td>Finish performing exclusive or non-exclusive on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td>pg_is_in_backup()</td>
<td>bool</td>
<td>True if an on-line exclusive backup is still in progress.</td>
</tr>
<tr>
<td>pg_backup_start_time()</td>
<td>timestamp with timezone</td>
<td>Get start time of an on-line exclusive backup in progress.</td>
</tr>
<tr>
<td>pg_switch_wal()</td>
<td>pg_lsn</td>
<td>Force switch to a new write-ahead log file (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td>pg_walfile_name(lsn pg_lsn)</td>
<td>text</td>
<td>Convert write-ahead log location to file name</td>
</tr>
<tr>
<td>pg_walfile_name_offset(lsn pg_lsn)</td>
<td>text, integer</td>
<td>Convert write-ahead log location to file name and decimal byte offset within file</td>
</tr>
<tr>
<td>pg_wal_lsn_diff(lsn pg_lsn, lsn pg_lsn)</td>
<td>numeric</td>
<td>Calculate the difference between two write-ahead log locations</td>
</tr>
</tbody>
</table>

pg_start_backup accepts an arbitrary user-defined label for the backup. (Typically this would be the name under which the backup dump file will be stored.) When used in exclusive mode, the function writes a backup label file (backup_label) and, if there are any links in the pg_tblspc directory, a tablespace map file (tablespace_map) into the database cluster's data directory, performs a checkpoint, and then returns the backup's starting write-ahead log location as text. The user can ignore this result value, but it is provided in case it is useful. When used in non-exclusive mode,
the contents of these files are instead returned by the `pg_stop_backup` function, and should be written to the backup by the caller.

```
postgres=# select pg_start_backup('label_goes_here');
pg_start_backup
-------------
0/D4445B8
(1 row)
```

There is an optional second parameter of type `boolean`. If `true`, it specifies executing `pg_start_backup` as quickly as possible. This forces an immediate checkpoint which will cause a spike in I/O operations, slowing any concurrently executing queries.

In an exclusive backup, `pg_stop_backup` removes the label file and, if it exists, the `tablespace_map` file created by `pg_start_backup`. In a non-exclusive backup, the contents of the `backup_label` and `tablespace_map` are returned in the result of the function, and should be written to files in the backup (and not in the data directory). There is an optional second parameter of type `boolean`. If false, the `pg_stop_backup` will return immediately after the backup is completed without waiting for WAL to be archived. This behavior is only useful for backup software which independently monitors WAL archiving. Otherwise, WAL required to make the backup consistent might be missing and make the backup useless. When this parameter is set to true, `pg_stop_backup` will wait for WAL to be archived when archiving is enabled; on the standby, this means that it will wait only when `archive_mode = always`. If write activity on the primary is low, it may be useful to run `pg_switch_wal` on the primary in order to trigger an immediate segment switch.

When executed on a primary, the function also creates a backup history file in the write-ahead log archive area. The history file includes the label given to `pg_start_backup`, the starting and ending write-ahead log locations for the backup, and the starting and ending times of the backup. The return value is the backup's ending write-ahead log location (which again can be ignored). After recording the ending location, the current write-ahead log insertion point is automatically advanced to the next write-ahead log file, so that the ending write-ahead log file can be archived immediately to complete the backup.

`pg_switch_wal` moves to the next write-ahead log file, allowing the current file to be archived (assuming you are using continuous archiving). The return value is the ending write-ahead log location + 1 within the just-completed write-ahead log file. If there has been no write-ahead log activity since the last write-ahead log switch, `pg_switch_wal` does nothing and returns the start location of the write-ahead log file currently in use.

`pg_create_restore_point` creates a named write-ahead log record that can be used as recovery target, and returns the corresponding write-ahead log location. The given name can then be used with `recovery_target_name` to specify the point up to which recovery will proceed. Avoid creating multiple restore points with the same name, since recovery will stop at the first one whose name matches the recovery target.

`pg_current_wal_lsn` displays the current write-ahead log write location in the same format used by the above functions. Similarly, `pg_current_wal_insert_lsn` displays the current write-ahead log insertion location and `pg_current_wal_flush_lsn` displays the current write-ahead log flush location. The insertion location is the “logical” end of the write-ahead log at any instant, while the write location is the end of what has actually been written out from the server's internal buffers and flush location is the location guaranteed to be written to durable storage. The write location is the end of what can be examined from outside the server, and is usually what you want if you are interested in archiving partially-complete write-ahead log files. The insertion and flush locations are made available primarily for server debugging purposes. These are both read-only operations and do not require superuser permissions.

You can use `pg_walfile_name_offset` to extract the corresponding write-ahead log file name and byte offset from the results of any of the above functions. For example:
Functions and Operators

postgres=# SELECT * FROM pg_walfile_name_offset(pg_stop_backup());
           file_name       | file_offset
--------------------------+-------------
00000001000000000000000D |     4039624
(1 row)

Similarly, pg_walfile_name extracts just the write-ahead log file name. When the given write-ahead log location is exactly at a write-ahead log file boundary, both these functions return the name of the preceding write-ahead log file. This is usually the desired behavior for managing write-ahead log archiving behavior, since the preceding file is the last one that currently needs to be archived.

pg_wal_lsn_diff calculates the difference in bytes between two write-ahead log locations. It can be used with pg_stat_replication or some functions shown in Table 9.79 to get the replication lag.

For details about proper usage of these functions, see Section 25.3.

9.26.4. Recovery Control Functions

The functions shown in Table 9.80 provide information about the current status of the standby. These functions may be executed both during recovery and in normal running.

Table 9.80. Recovery Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_is_in_recovery()</td>
<td>bool</td>
<td>True if recovery is still in progress.</td>
</tr>
<tr>
<td>pg_last_wal_receive_lsn()</td>
<td>pg_lsn</td>
<td>Get last write-ahead log location received and synced to disk by streaming replication. While streaming replication is in progress this will increase monotonically. If recovery has completed this will remain static at the value of the last WAL record received and synced to disk during recovery. If streaming replication is disabled, or if it has not yet started, the function returns NULL.</td>
</tr>
<tr>
<td>pg_last_wal_replay_lsn()</td>
<td>pg_lsn</td>
<td>Get last write-ahead log location replayed during recovery. If recovery is still in progress this will increase monotonically. If recovery has completed then this value will remain static at the value of the last WAL record applied during that recovery. When the server has been started normally without recovery the function returns NULL.</td>
</tr>
<tr>
<td>pg_last_xact_replay_timestamp()</td>
<td>timestamp with timezone</td>
<td>Get time stamp of last transaction replayed during recovery. This is the time at which the commit or abort WAL record for that transaction was generated on the primary. If no trans-</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_is_wal_replayPaused()</td>
<td>bool</td>
<td>True if recovery is paused.</td>
</tr>
<tr>
<td>pg_wal_replayPause()</td>
<td>void</td>
<td>Pauses recovery immediately (restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
<tr>
<td>pg_wal_replayResume()</td>
<td>void</td>
<td>Restarts recovery if it was paused (restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
</tbody>
</table>

The functions shown in Table 9.81 control the progress of recovery. These functions may be executed only during recovery.

#### Table 9.81. Recovery Control Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_is_wal_replayPaused()</td>
<td>bool</td>
<td>True if recovery is paused.</td>
</tr>
<tr>
<td>pg_wal_replayPause()</td>
<td>void</td>
<td>Pauses recovery immediately (restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
<tr>
<td>pg_wal_replayResume()</td>
<td>void</td>
<td>Restarts recovery if it was paused (restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
</tbody>
</table>

While recovery is paused no further database changes are applied. If in hot standby, all new queries will see the same consistent snapshot of the database, and no further query conflicts will be generated until recovery is resumed.

If streaming replication is disabled, the paused state may continue indefinitely without problem. While streaming replication is in progress WAL records will continue to be received, which will eventually fill available disk space, depending upon the duration of the pause, the rate of WAL generation and available disk space.

### 9.26.5. Snapshot Synchronization Functions

PostgreSQL allows database sessions to synchronize their snapshots. A snapshot determines which data is visible to the transaction that is using the snapshot. Synchronized snapshots are necessary when two or more sessions need to see identical content in the database. If two sessions just start their transactions independently, there is always a possibility that some third transaction commits between the executions of the two `START TRANSACTION` commands, so that one session sees the effects of that transaction and the other does not.

To solve this problem, PostgreSQL allows a transaction to `export` the snapshot it is using. As long as the exporting transaction remains open, other transactions can `import` its snapshot, and thereby be guaranteed that they see exactly the same view of the database that the first transaction sees. But note that any database changes made by any one of these transactions remain invisible to the other transactions, as is usual for changes made by uncommitted transactions. So the transactions are synchronized with respect to pre-existing data, but act normally for changes they make themselves.
Snapshots are exported with the `pg_export_snapshot` function, shown in Table 9.82, and imported with the SET TRANSACTION command.

### Table 9.82. Snapshot Synchronization Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_export_snapshot()</code></td>
<td>text</td>
<td>Save the current snapshot and return its identifier</td>
</tr>
</tbody>
</table>

The function `pg_export_snapshot` saves the current snapshot and returns a `text` string identifying the snapshot. This string must be passed (outside the database) to clients that want to import the snapshot. The snapshot is available for import only until the end of the transaction that exported it. A transaction can export more than one snapshot, if needed. Note that doing so is only useful in READ COMMITTED transactions, since in REPEATABLE READ and higher isolation levels, transactions use the same snapshot throughout their lifetime. Once a transaction has exported any snapshots, it cannot be prepared with PREPARE TRANSACTION.

See SET TRANSACTION for details of how to use an exported snapshot.

### 9.26.6. Replication Functions

The functions shown in Table 9.83 are for controlling and interacting with replication features. See Section 26.2.5, Section 26.2.6, and Chapter 50 for information about the underlying features. Use of these functions is restricted to superusers.

Many of these functions have equivalent commands in the replication protocol; see Section 53.4.

The functions described in Section 9.26.3, Section 9.26.4, and Section 9.26.5 are also relevant for replication.

### Table 9.83. Replication SQL Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_create_physical_replication_slot(slot_name name [, immediately_reserve boolean, temporary boolean])</code></td>
<td><code>(slot_name name, lsn pg_lsn)</code></td>
<td>Creates a new physical replication slot named <code>slot_name</code>. The optional second parameter, when true, specifies that the LSN for this replication slot be reserved immediately; otherwise the LSN is reserved on first connection from a streaming replication client. Streaming changes from a physical slot is only possible with the streaming-replication protocol — see Section 53.4. The optional third parameter, <code>temporary</code>, when set to true, specifies that the slot should not be permanently stored to disk and is only meant for use by current session. Temporary slots are also released upon any error. This function corresponds to the replication protocol command <code>CREATE_REPLICATION_SLOT ... PHYSICAL</code>.</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
</tr>
<tr>
<td>pg_drop_replication_slot(slot_name name)</td>
<td>void</td>
<td>Drops the physical or logical replication slot named <code>slot_name</code>. Same as replication protocol command <code>DROP_REPLICATION_SLOT</code>. For logical slots, this must be called when connected to the same database the slot was created on.</td>
</tr>
<tr>
<td>pg_create_logical_replication_slot(slot_name name, plugin name [, temporary boolean])</td>
<td>(slot_name name, lsn pg_lsn)</td>
<td>Creates a new logical (decoding) replication slot named <code>slot_name</code> using the output plugin <code>plugin</code>. The optional third parameter, <code>temporary</code>, when set to true, specifies that the slot should not be permanently stored to disk and is only meant for use by current session. Temporary slots are also released upon any error. A call to this function has the same effect as the replication protocol command <code>CREATE_REPLICATION_SLOT ... LOGICAL</code>.</td>
</tr>
<tr>
<td>pg_logical_slot_get_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</td>
<td>(lsn pg_lsn, xid xid, data text)</td>
<td>Returns changes in the slot <code>slot_name</code>, starting from the point at which since changes have been consumed last. If <code>upto_lsn</code> and <code>upto_nchanges</code> are NULL, logical decoding will continue until end of WAL. If <code>upto_lsn</code> is non-NULL, decoding will include only those transactions which commit prior to the specified LSN. If <code>upto_nchanges</code> is non-NULL, decoding will stop when the number of rows produced by decoding exceeds the specified value. Note, however, that the actual number of rows returned may be larger, since this limit is only checked after adding the rows produced when decoding each new transaction commit.</td>
</tr>
<tr>
<td>pg_logical_slot_peek_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</td>
<td>(lsn pg_lsn, xid xid, data text)</td>
<td>Behaves just like the <code>pg_logical_slot_get_changes()</code> function, except that changes are not consumed; that is, they will be returned again on future calls.</td>
</tr>
<tr>
<td>pg_logical_slot_get_binary(slot_lsn pg_lsn, xid xid, data bytea)</td>
<td></td>
<td>Behaves just like the <code>pg_logical_slot_get_changes()</code> function, except that changes are not consumed; that is, they will be returned again on future calls.</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------------------------</td>
<td>-------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>ry_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</td>
<td></td>
<td>Returns the changes as bytea, unlike <code>lot_get_changes()</code>, which returns them as a function.</td>
</tr>
<tr>
<td>pg_logical_slot_peek_binary_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</td>
<td>(lsn pg_lsn, xid xid, data bytea)</td>
<td>Behaves just like the <code>pg_logical_slot_get_changes()</code> function, except that changes are returned as bytea and that changes are not consumed; that is, they will be returned again on future calls.</td>
</tr>
<tr>
<td>pg_replication_slot_advance(slot_name name, upto_lsn pg_lsn)</td>
<td>(slot_name name, end lsn pg_lsn) bool</td>
<td>Advances the current confirmed position of a replication slot named <code>slot_name</code>. The slot will not be moved backwards, and it will not be moved beyond the current insert location. Returns name of the slot and real position to which it was advanced to.</td>
</tr>
<tr>
<td>pg_replication_origin_create(node_name text)</td>
<td>oid</td>
<td>Creates a replication origin with the given external name, and return the internal id assigned to it.</td>
</tr>
<tr>
<td>pg_replication_origin_drop(node_name text)</td>
<td>void</td>
<td>Deletes a previously created replication origin, including any associated replay progress.</td>
</tr>
<tr>
<td>pg_replication_origin_oid(node_name text)</td>
<td>oid</td>
<td>Lookup a replication origin by name and return the internal id. If no corresponding replication origin is found an error is thrown.</td>
</tr>
<tr>
<td>pg_replication_origin_session_setup(node_name text)</td>
<td>void</td>
<td>Mark the current session as replaying from the given origin, allowing replay progress to be tracked. Use <code>pg_replication_origin_session_reset</code> to revert. Can only be used if no previous origin is configured.</td>
</tr>
<tr>
<td>pg_replication_origin_session_reset()</td>
<td>void</td>
<td>Cancel the effects of <code>pg_replication_origin_session_setup()</code>.</td>
</tr>
<tr>
<td>pg_replication_origin_session_is_set_up()</td>
<td>bool</td>
<td>Has a replication origin been configured in the current session?</td>
</tr>
<tr>
<td>pg_replication_origin_session_progress(flush bool)</td>
<td>pg_lsn</td>
<td>Return the replay location for the replication origin configured in the current session. The parameter <code>flush</code> determines whether the corresponding local transaction will be guaranteed to have been flushed to disk or not.</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------------------------------------</td>
<td>-------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_replication_origin_xact_setup(ori-gin_lsn pg_lsn, ori-gin_timestamp time-stamptz)</code></td>
<td>void</td>
<td>Mark the current transaction as replaying a transaction that has committed at the given LSN and timestamp. Can only be called when a replication origin has previously been configured using <code>pg_replication_origin_session_setup()</code>.</td>
</tr>
<tr>
<td><code>pg_replication_origin_xact_reset()</code></td>
<td>void</td>
<td>Cancel the effects of <code>pg_replication_origin_xact_setup()</code>.</td>
</tr>
<tr>
<td><code>pg_replication_origin_advance(node_name text, lsn pg_lsn)</code></td>
<td>void</td>
<td>Set replication progress for the given node to the given location.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>This primarily is useful for setting up the initial location or a new location after configuration changes and similar. Be aware that careless use of this function can lead to inconsistently replicated data.</td>
</tr>
<tr>
<td><code>pg_replication_origin_progress(node_name text, flush bool)</code></td>
<td>pg_lsn</td>
<td>Return the replay location for the given replication origin. The parameter <code>flush</code> determines whether the corresponding local transaction will be guaranteed to have been flushed to disk or not.</td>
</tr>
<tr>
<td><code>pg_logical_emit_message(transactional bool, prefix text, content text)</code></td>
<td>pg_lsn</td>
<td>Emit text logical decoding message. This can be used to pass generic messages to logical decoding plugins through WAL. The parameter <code>transactional</code> specifies if the message should be part of current transaction or if it should be written immediately and decoded as soon as the logical decoding reads the record. The <code>prefix</code> is textual prefix used by the logical decoding plugins to easily recognize interesting messages for them. The <code>content</code> is the text of the message.</td>
</tr>
</tbody>
</table>
| `pg_logical_emit_message(transactional bool, prefix text, content bytea)` | pg_lsn      | Emit binary logical decoding message. This can be used to pass generic messages to logical decoding plugins through WAL. The parameter `transactional` specifies if the message should be part of current transaction or if it should be written immediately and decoded as soon as the logical decoding reads the record. The `prefix` is textual prefix used by
9.26.7. Database Object Management Functions

The functions shown in Table 9.84 calculate the disk space usage of database objects.

Table 9.84. Database Object Size Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_column_size(any)</td>
<td>int</td>
<td>Number of bytes used to store a particular value (possibly compressed)</td>
</tr>
<tr>
<td>pg_database_size(oid)</td>
<td>bigint</td>
<td>Disk space used by the database with the specified OID</td>
</tr>
<tr>
<td>pg_database_size(name)</td>
<td>bigint</td>
<td>Disk space used by the database with the specified name</td>
</tr>
<tr>
<td>pg_indexes_size(regclass)</td>
<td>bigint</td>
<td>Total disk space used by indexes attached to the specified table</td>
</tr>
<tr>
<td>pg_relation_size(regclass, fork text)</td>
<td>bigint</td>
<td>Disk space used by the specified fork ('main', '_fsm', 'vm', or 'init') of the specified table or index</td>
</tr>
<tr>
<td>pg_relation_size(regclass)</td>
<td>bigint</td>
<td>Shorthand for pg_relationsize(..., 'main')</td>
</tr>
<tr>
<td>pg_size_bytes(text)</td>
<td>bigint</td>
<td>Converts a size in human-readable format with size units into bytes</td>
</tr>
<tr>
<td>pg_size_pretty(bigint)</td>
<td>text</td>
<td>Converts a size in bytes expressed as a 64-bit integer into a human-readable format with size units</td>
</tr>
<tr>
<td>pg_size_pretty(numeric)</td>
<td>text</td>
<td>Converts a size in bytes expressed as a numeric value into a human-readable format with size units</td>
</tr>
<tr>
<td>pg_table_size(regclass)</td>
<td>bigint</td>
<td>Disk space used by the specified table, excluding indexes (but including TOAST, free space map, and visibility map)</td>
</tr>
<tr>
<td>pg_tablespace_size(oid)</td>
<td>bigint</td>
<td>Disk space used by the tablespace with the specified OID</td>
</tr>
<tr>
<td>pg_tablespace_size(name)</td>
<td>bigint</td>
<td>Disk space used by the tablespace with the specified name</td>
</tr>
<tr>
<td>pg_total_relation_size(regclass)</td>
<td>bigint</td>
<td>Total disk space used by the specified table, including all indexes and TOAST data</td>
</tr>
</tbody>
</table>
pg_column_size shows the space used to store any individual data value.

pg_total_relation_size accepts the OID or name of a table or toast table, and returns the total on-disk space used for that table, including all associated indexes. This function is equivalent to pg_table_size + pg_indexes_size.

pg_table_size accepts the OID or name of a table and returns the disk space needed for that table, exclusive of indexes. (TOAST space, free space map, and visibility map are included.)

pg_indexes_size accepts the OID or name of a table and returns the total disk space used by all the indexes attached to that table.

pg_database_size and pg_tablespace_size accept the OID or name of a database or tablespace, and return the total disk space used therein. To use pg_database_size, you must have CONNECT permission on the specified database (which is granted by default), or be a member of the pg_read_all_stats role. To use pg_tablespace_size, you must have CREATE permission on the specified tablespace, or be a member of the pg_read_all_stats role unless it is the default tablespace for the current database.

pg_relation_size accepts the OID or name of a table, index or toast table, and returns the on-disk size in bytes of one fork of that relation. (Note that for most purposes it is more convenient to use the higher-level functions pg_total_relation_size or pg_table_size, which sum the sizes of all forks.) With one argument, it returns the size of the main data fork of the relation. The second argument can be provided to specify which fork to examine:

- 'main' returns the size of the main data fork of the relation.
- 'fsm' returns the size of the Free Space Map (see Section 68.3) associated with the relation.
- 'vm' returns the size of the Visibility Map (see Section 68.4) associated with the relation.
- 'init' returns the size of the initialization fork, if any, associated with the relation.

pg_size_pretty can be used to format the result of one of the other functions in a human-readable way, using bytes, kB, MB, GB or TB as appropriate.

pg_size_bytes can be used to get the size in bytes from a string in human-readable format. The input may have units of bytes, kB, MB, GB or TB, and is parsed case-insensitively. If no units are specified, bytes are assumed.

Note

The units kB, MB, GB and TB used by the functions pg_size_pretty and pg_size_bytes are defined using powers of 2 rather than powers of 10, so 1kB is 1024 bytes, 1MB is $1024^2 = 1048576$ bytes, and so on.

The functions above that operate on tables or indexes accept a regclass argument, which is simply the OID of the table or index in the pg_class system catalog. You do not have to look up the OID by hand, however, since the regclass data type's input converter will do the work for you. Just write the table name enclosed in single quotes so that it looks like a literal constant. For compatibility with the handling of ordinary SQL names, the string will be converted to lower case unless it contains double quotes around the table name.

If an OID that does not represent an existing object is passed as argument to one of the above functions, NULL is returned.

The functions shown in Table 9.85 assist in identifying the specific disk files associated with database objects.
### Table 9.85. Database Object Location Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_relation_filenode(relation regclass)</code></td>
<td>oid</td>
<td>Filenode number of the specified relation</td>
</tr>
<tr>
<td><code>pg_relation_filepath(relation regclass)</code></td>
<td>text</td>
<td>File path name of the specified relation</td>
</tr>
<tr>
<td><code>pg_filenode_relation(tablespace oid, filenode oid)</code></td>
<td>regclass</td>
<td>Find the relation associated with a given tablespace and filenode</td>
</tr>
</tbody>
</table>

`pg_relation_filenode` accepts the OID or name of a table, index, sequence, or toast table, and returns the “filenode” number currently assigned to it. The filenode is the base component of the file name(s) used for the relation (see Section 68.1 for more information). For most tables the result is the same as `pg_class.relfilenode`, but for certain system catalogs `relfilenode` is zero and this function must be used to get the correct value. The function returns NULL if passed a relation that does not have storage, such as a view.

`pg_relation_filepath` is similar to `pg_relation_filenode`, but it returns the entire file path name (relative to the database cluster’s data directory `PGDATA`) of the relation.

`pg_filenode_relation` is the reverse of `pg_relation_filenode`. Given a “tablespace” OID and a “filenode”, it returns the associated relation’s OID. For a table in the database’s default tablespace, the tablespace can be specified as 0.

Table 9.86 lists functions used to manage collations.

### Table 9.86. Collation Management Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_collation_actual_version(oid)</code></td>
<td>text</td>
<td>Return actual version of collation from operating system</td>
</tr>
<tr>
<td><code>pg_import_system_collations(schema regnamespace)</code></td>
<td>integer</td>
<td>Import operating system collations</td>
</tr>
</tbody>
</table>

`pg_collation_actual_version` returns the actual version of the collation object as it is currently installed in the operating system. If this is different from the value in `pg_collation.relversion`, then objects depending on the collation might need to be rebuilt. See also `ALTER COLLATION`.

`pg_import_system_collations` adds collations to the system catalog `pg_collation` based on all the locales it finds in the operating system. This is what `initdb` uses; see Section 23.2.2 for more details. If additional locales are installed into the operating system later on, this function can be run again to add collations for the new locales. Locales that match existing entries in `pg_collation` will be skipped. (But collation objects based on locales that are no longer present in the operating system are not removed by this function.) The `schema` parameter would typically be `pg_catalog`, but that is not a requirement; the collations could be installed into some other schema as well. The function returns the number of new collation objects it created.

### 9.26.8. Index Maintenance Functions

Table 9.87 shows the functions available for index maintenance tasks. These functions cannot be executed during recovery. Use of these functions is restricted to superusers and the owner of the given index.
### Table 9.87. Index Maintenance Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>brin_summarize_new_values(index regclass)</td>
<td>integer</td>
<td>summarize page ranges not already summarized</td>
</tr>
<tr>
<td>brin_summarize_range(index regclass, blockNumber bigint)</td>
<td>integer</td>
<td>summarize the page range covering the given block, if not already summarized</td>
</tr>
<tr>
<td>brin_desummarize_range(index regclass, blockNumber bigint)</td>
<td>integer</td>
<td>de-summarize the page range covering the given block, if summarized</td>
</tr>
<tr>
<td>gin_clean_pending_list(index regclass)</td>
<td>bigint</td>
<td>move GIN pending list entries into main index structure</td>
</tr>
</tbody>
</table>

brin_summarize_new_values accepts the OID or name of a BRIN index and inspects the index to find page ranges in the base table that are not currently summarized by the index; for any such range it creates a new summary index tuple by scanning the table pages. It returns the number of new page range summaries that were inserted into the index. brin_summarize_range does the same, except it only summarizes the range that covers the given block number.

gin_clean_pending_list accepts the OID or name of a GIN index and cleans up the pending list of the specified index by moving entries in it to the main GIN data structure in bulk. It returns the number of pages removed from the pending list. Note that if the argument is a GIN index built with the fastupdate option disabled, no cleanup happens and the return value is 0, because the index doesn't have a pending list. Please see Section 66.4.1 and Section 66.5 for details of the pending list and fastupdate option.


The functions shown in Table 9.88 provide native access to files on the machine hosting the server. Only files within the database cluster directory and the log_directory can be accessed unless the user is granted the role pg_read_server_files. Use a relative path for files in the cluster directory, and a path matching the log_directory configuration setting for log files.

Note that granting users the EXECUTE privilege on the pg_read_file(), or related, functions allows them the ability to read any file on the server which the database can read and that those reads bypass all in-database privilege checks. This means that, among other things, a user with this access is able to read the contents of the pg_authid table where authentication information is contained, as well as read any file in the database. Therefore, granting access to these functions should be carefully considered.

### Table 9.88. Generic File Access Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_ls_dir(dirname text [, missing_ok boolean, include_dot_dirs boolean])</td>
<td>setof text</td>
<td>List the contents of a directory. Restricted to superusers by default, but other users can be granted EXECUTE to run the function.</td>
</tr>
<tr>
<td>pg_ls_logdir()</td>
<td>setof record</td>
<td>List the name, size, and last modification time of files in the log</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-------------------</td>
<td>----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>pg_ls_waldir()</td>
<td>setof record</td>
<td>List the name, size, and last modification time of files in the WAL directory. Access is granted to members of the pg_monitor role and may be granted to other non-superuser roles.</td>
</tr>
<tr>
<td>pg_read_file(filename text [, offset bigint, length bigint [, missing_ok boolean] ]))</td>
<td>text</td>
<td>Return the contents of a text file. Restricted to superusers by default, but other users can be granted EXECUTE to run the function.</td>
</tr>
<tr>
<td>pg_read_binary_file(filename text [, offset bigint, length bigint [, missing_ok boolean] ]))</td>
<td>bytea</td>
<td>Return the contents of a file. Restricted to superusers by default, but other users can be granted EXECUTE to run the function.</td>
</tr>
<tr>
<td>pg_stat_file(filename text [, missing_ok boolean])</td>
<td>record</td>
<td>Return information about a file. Restricted to superusers by default, but other users can be granted EXECUTE to run the function.</td>
</tr>
</tbody>
</table>

Some of these functions take an optional `missing_ok` parameter, which specifies the behavior when the file or directory does not exist. If `true`, the function returns NULL (except `pg_ls_dir`, which returns an empty result set). If `false`, an error is raised. The default is `false`.

- **pg_ls_dir** returns the names of all files (and directories and other special files) in the specified directory. The `include_dot_dirs` indicates whether `.` and `..` are included in the result set. The default is to exclude them (false), but including them can be useful when `missing_ok` is true, to distinguish an empty directory from a non-existent directory.

- **pg_ls_logdir** returns the name, size, and last modified time (mtime) of each file in the log directory. By default, only superusers and members of the `pg_monitor` role can use this function. Access may be granted to others using GRANT.

- **pg_ls_waldir** returns the name, size, and last modified time (mtime) of each file in the write ahead log (WAL) directory. By default only superusers and members of the `pg_monitor` role can use this function. Access may be granted to others using GRANT.

- **pg_read_file** returns part of a text file, starting at the given `offset`, returning at most `length` bytes (less if the end of file is reached first). If `offset` is negative, it is relative to the end of the file. If `offset` and `length` are omitted, the entire file is returned. The bytes read from the file are interpreted as a string in the server encoding; an error is thrown if they are not valid in that encoding.

- **pg_read_binary_file** is similar to `pg_read_file`, except that the result is a `bytea` value; accordingly, no encoding checks are performed. In combination with the `convert_from` function, this function can be used to read a file in a specified encoding.
```sql
SELECT convert_from(pg_read_binary_file('file_in_utf8.txt'), 'UTF8');
```

`pg_stat_file` returns a record containing the file size, last accessed time stamp, last modified time stamp, last file status change time stamp (Unix platforms only), file creation time stamp (Windows only), and a boolean indicating if it is a directory. Typical usages include:

```sql
SELECT * FROM pg_stat_file('filename');
SELECT (pg_stat_file('filename')).modification;
```

### 9.26.10. Advisory Lock Functions

The functions shown in Table 9.89 manage advisory locks. For details about proper use of these functions, see Section 13.3.5.

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_advisory_lock(key bigint)</code></td>
<td>void</td>
<td>Obtain exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_lock(key1 int, key2 int)</code></td>
<td>void</td>
<td>Obtain exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_lock共享(key bigint)</code></td>
<td>void</td>
<td>Obtain shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_lock_shared(key1 int, key2 int)</code></td>
<td>void</td>
<td>Obtain shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock(key bigint)</code></td>
<td>boolean</td>
<td>Release an exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Release an exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_all()</code></td>
<td>void</td>
<td>Release all session level advisory locks held by the current session</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_shared(key bigint)</code></td>
<td>boolean</td>
<td>Release a shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_shared(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Release a shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock(key bigint)</code></td>
<td>void</td>
<td>Obtain exclusive transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock(key1 int, key2 int)</code></td>
<td>void</td>
<td>Obtain exclusive transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock共享(key bigint)</code></td>
<td>void</td>
<td>Obtain shared transaction level advisory lock</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------------------</td>
<td>-------------</td>
<td>-------------------------------------------------------</td>
</tr>
<tr>
<td>pg_advisory_xact_lock_shared(key1 int, key2 int)</td>
<td>void</td>
<td>Obtain shared transaction level advisory lock</td>
</tr>
<tr>
<td>pg_try_advisory_lock(key bigint)</td>
<td>boolean</td>
<td>Obtain exclusive session level advisory lock if available</td>
</tr>
<tr>
<td>pg_try_advisory_lock(key1 int, key2 int)</td>
<td>boolean</td>
<td>Obtain exclusive session level advisory lock if available</td>
</tr>
<tr>
<td>pg_try_advisory_lock_shared(key bigint)</td>
<td>boolean</td>
<td>Obtain shared session level advisory lock if available</td>
</tr>
<tr>
<td>pg_try_advisory_lock_shared(key1 int, key2 int)</td>
<td>boolean</td>
<td>Obtain shared session level advisory lock if available</td>
</tr>
<tr>
<td>pg_advisory_unlock</td>
<td>boolean</td>
<td>Release a previously-acquired exclusive lock</td>
</tr>
<tr>
<td>pg_advisory_unlock_shared</td>
<td>boolean</td>
<td>Release a previously-acquired shared lock</td>
</tr>
</tbody>
</table>

pg_advisory_lock locks an application-defined resource, which can be identified either by a single 64-bit key value or two 32-bit key values (note that these two key spaces do not overlap). If another session already holds a lock on the same resource identifier, this function will wait until the resource becomes available. The lock is exclusive. Multiple lock requests stack, so that if the same resource is locked three times it must then be unlocked three times to be released for other sessions' use.

pg_advisory_lock_shared works the same as pg_advisory_lock, except the lock can be shared with other sessions requesting shared locks. Only would-be exclusive lockers are locked out.

pg_try_advisory_lock is similar to pg_advisory_lock, except the function will not wait for the lock to become available. It will either obtain the lock immediately and return true, or return false if the lock cannot be acquired immediately.

pg_try_advisory_lock_shared works the same as pg_try_advisory_lock, except it attempts to acquire a shared rather than an exclusive lock.

pg_advisory_unlock will release a previously-acquired exclusive session level advisory lock. It returns true if the lock is successfully released. If the lock was not held, it will return false, and in addition, an SQL warning will be reported by the server.

pg_advisory_unlock_shared works the same as pg_advisory_unlock, except it releases a shared session level advisory lock.
pg_advisory_unlock_all will release all session level advisory locks held by the current session. (This function is implicitly invoked at session end, even if the client disconnects ungracefully.)

pg_advisory_xact_lock works the same as pg_advisory_lock, except the lock is automatically released at the end of the current transaction and cannot be released explicitly.

pg_advisory_xact_lock_shared works the same as pg_advisory_lock_shared, except the lock is automatically released at the end of the current transaction and cannot be released explicitly.

pg_try_advisory_xact_lock works the same as pg_try_advisory_lock, except the lock, if acquired, is automatically released at the end of the current transaction and cannot be released explicitly.

pg_try_advisory_xact_lock_shared works the same as pg_try_advisory_lock_shared, except the lock, if acquired, is automatically released at the end of the current transaction and cannot be released explicitly.

9.27. Trigger Functions

Currently PostgreSQL provides one built in trigger function, suppress_redundant_updates_trigger, which will prevent any update that does not actually change the data in the row from taking place, in contrast to the normal behavior which always performs the update regardless of whether or not the data has changed. (This normal behavior makes updates run faster, since no checking is required, and is also useful in certain cases.)

Ideally, you should normally avoid running updates that don't actually change the data in the record. Redundant updates can cost considerable unnecessary time, especially if there are lots of indexes to alter, and space in dead rows that will eventually have to be vacuumed. However, detecting such situations in client code is not always easy, or even possible, and writing expressions to detect them can be error-prone. An alternative is to use suppress_redundant_updates_trigger, which will skip updates that don't change the data. You should use this with care, however. The trigger takes a small but non-trivial time for each record, so if most of the records affected by an update are actually changed, use of this trigger will actually make the update run slower.

The suppress_redundant_updates_trigger function can be added to a table like this:

```sql
CREATE TRIGGER z_min_update
BEFORE UPDATE ON tablename
FOR EACH ROW EXECUTE FUNCTION suppress_redundant_updates_trigger();
```

In most cases, you would want to fire this trigger last for each row. Bearing in mind that triggers fire in name order, you would then choose a trigger name that comes after the name of any other trigger you might have on the table.

For more information about creating triggers, see CREATE TRIGGER.

9.28. Event Trigger Functions

PostgreSQL provides these helper functions to retrieve information from event triggers.

For more information about event triggers, see Chapter 40.

9.28.1. Capturing Changes at Command End
pg_event_trigger_ddl_commands returns a list of DDL commands executed by each user action, when invoked in a function attached to a ddl_command_end event trigger. If called in any other context, an error is raised. pg_event_trigger_ddl_commands returns one row for each base command executed; some commands that are a single SQL sentence may return more than one row. This function returns the following columns:

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>OID of catalog the object belongs in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>OID of the object itself</td>
</tr>
<tr>
<td>objsubid</td>
<td>integer</td>
<td>Sub-object ID (e.g. attribute number for a column)</td>
</tr>
<tr>
<td>command_tag</td>
<td>text</td>
<td>Command tag</td>
</tr>
<tr>
<td>object_type</td>
<td>text</td>
<td>Type of the object</td>
</tr>
<tr>
<td>schema_name</td>
<td>text</td>
<td>Name of the schema the object belongs in, if any; otherwise NULL. No quoting is applied.</td>
</tr>
<tr>
<td>object_identity</td>
<td>text</td>
<td>Text rendering of the object identity, schema-qualified. Each identifier included in the identity is quoted if necessary.</td>
</tr>
<tr>
<td>in_extension</td>
<td>bool</td>
<td>True if the command is part of an extension script</td>
</tr>
<tr>
<td>command</td>
<td>pg_ddl_command</td>
<td>A complete representation of the command, in internal format. This cannot be output directly, but it can be passed to other functions to obtain different pieces of information about the command.</td>
</tr>
</tbody>
</table>

9.28.2. Processing Objects Dropped by a DDL Command

pg_event_trigger_dropped_objects returns a list of all objects dropped by the command in whose sql_drop event it is called. If called in any other context, pg_event_trigger_dropped_objects raises an error. pg_event_trigger_dropped_objects returns the following columns:

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>OID of catalog the object belonged in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>OID of the object itself</td>
</tr>
<tr>
<td>objsubid</td>
<td>integer</td>
<td>Sub-object ID (e.g. attribute number for a column)</td>
</tr>
<tr>
<td>original</td>
<td>bool</td>
<td>True if this was one of the root object(s) of the deletion</td>
</tr>
<tr>
<td>normal</td>
<td>bool</td>
<td>True if there was a normal dependency relationship in the dependency graph leading to this object</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>is_temporary</td>
<td>bool</td>
<td>True if this was a temporary object</td>
</tr>
<tr>
<td>object_type</td>
<td>text</td>
<td>Type of the object</td>
</tr>
<tr>
<td>schema_name</td>
<td>text</td>
<td>Name of the schema the object belonged in, if any; otherwise NULL. No quoting is applied.</td>
</tr>
<tr>
<td>object_name</td>
<td>text</td>
<td>Name of the object, if the combination of schema and name can be used as a unique identifier for the object; otherwise NULL. No quoting is applied, and name is never schema-qualified.</td>
</tr>
<tr>
<td>object_identity</td>
<td>text</td>
<td>Text rendering of the object identity, schema-qualified. Each identifier included in the identity is quoted if necessary.</td>
</tr>
<tr>
<td>address_names</td>
<td>text[]</td>
<td>An array that, together with object_type and address_args, can be used by the pg_get_object_address() function to recreate the object address in a remote server containing an identically named object of the same kind.</td>
</tr>
<tr>
<td>address_args</td>
<td>text[]</td>
<td>Complement for address_names</td>
</tr>
</tbody>
</table>

The `pg_event_trigger_dropped_objects` function can be used in an event trigger like this:

```sql
CREATE FUNCTION test_event_trigger_for_drops()
RETURNS event_trigger LANGUAGE plpgsql AS $$
DECLARE
    obj record;
BEGIN
    FOR obj IN SELECT * FROM pg_event_trigger_dropped_objects()
    LOOP
        RAISE NOTICE '% dropped object: % %.% %',
        tg_tag, obj.object_type, obj.schema_name, obj.object_name, obj.object_identity;
    END LOOP;
END
$$;
CREATE EVENT TRIGGER test_event_trigger_for_drops
ON sql_drop
EXECUTE FUNCTION test_event_trigger_for_drops();
```

### 9.28.3. Handling a Table Rewrite Event

The functions shown in Table 9.90 provide information about a table for which a `table_rewrite` event has just been called. If called in any other context, an error is raised.
Table 9.90. Table Rewrite information

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_event_trigger_table_rewrite_oid()</td>
<td>Oid</td>
<td>The OID of the table about to be rewritten.</td>
</tr>
<tr>
<td>pg_event_trigger_table_rewrite_reason()</td>
<td>int</td>
<td>The reason code(s) explaining the reason for rewriting. The exact meaning of the codes is release dependent.</td>
</tr>
</tbody>
</table>

The `pg_event_trigger_table_rewrite_oid` function can be used in an event trigger like this:

```sql
CREATE FUNCTION test_event_trigger_table_rewrite_oid()
RETURNS event_trigger
LANGUAGE plpgsql AS
$$
BEGIN
    RAISE NOTICE 'rewriting table % for reason %',
    pg_event_trigger_table_rewrite_oid()::regclass,
    pg_event_trigger_table_rewrite_reason();
END;
$$;

CREATE EVENT TRIGGER test_table_rewrite_oid
ON table_rewrite
    EXECUTE FUNCTION test_event_trigger_table_rewrite_oid();
```

```sql
```
Chapter 10. Type Conversion

SQL statements can, intentionally or not, require the mixing of different data types in the same expression. PostgreSQL has extensive facilities for evaluating mixed-type expressions.

In many cases a user does not need to understand the details of the type conversion mechanism. However, implicit conversions done by PostgreSQL can affect the results of a query. When necessary, these results can be tailored by using explicit type conversion.

This chapter introduces the PostgreSQL type conversion mechanisms and conventions. Refer to the relevant sections in Chapter 8 and Chapter 9 for more information on specific data types and allowed functions and operators.

10.1. Overview

SQL is a strongly typed language. That is, every data item has an associated data type which determines its behavior and allowed usage. PostgreSQL has an extensible type system that is more general and flexible than other SQL implementations. Hence, most type conversion behavior in PostgreSQL is governed by general rules rather than by ad hoc heuristics. This allows the use of mixed-type expressions even with user-defined types.

The PostgreSQL scanner/parser divides lexical elements into five fundamental categories: integers, non-integer numbers, strings, identifiers, and key words. Constants of most non-numeric types are first classified as strings. The SQL language definition allows specifying type names with strings, and this mechanism can be used in PostgreSQL to start the parser down the correct path. For example, the query:

```
SELECT text 'Origin' AS "label", point '(0,0)' AS "value";
```

<table>
<thead>
<tr>
<th>label</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Origin</td>
<td>(0,0)</td>
</tr>
</tbody>
</table>

(1 row)

has two literal constants, of type text and point. If a type is not specified for a string literal, then the placeholder type unknown is assigned initially, to be resolved in later stages as described below.

There are four fundamental SQL constructs requiring distinct type conversion rules in the PostgreSQL parser:

Function calls

Much of the PostgreSQL type system is built around a rich set of functions. Functions can have one or more arguments. Since PostgreSQL permits function overloading, the function name alone does not uniquely identify the function to be called; the parser must select the right function based on the data types of the supplied arguments.

Operators

PostgreSQL allows expressions with prefix and postfix unary (one-argument) operators, as well as binary (two-argument) operators. Like functions, operators can be overloaded, so the same problem of selecting the right operator exists.

Value Storage

SQL INSERT and UPDATE statements place the results of expressions into a table. The expressions in the statement must be matched up with, and perhaps converted to, the types of the target columns.
Type Conversion

UNION, CASE, and related constructs

Since all query results from a unionized SELECT statement must appear in a single set of columns, the types of the results of each SELECT clause must be matched up and converted to a uniform set. Similarly, the result expressions of a CASE construct must be converted to a common type so that the CASE expression as a whole has a known output type. The same holds for ARRAY constructs, and for the GREATEST and LEAST functions.

The system catalogs store information about which conversions, or casts, exist between which data types, and how to perform those conversions. Additional casts can be added by the user with the CREATE CAST command. (This is usually done in conjunction with defining new data types. The set of casts between built-in types has been carefully crafted and is best not altered.)

An additional heuristic provided by the parser allows improved determination of the proper casting behavior among groups of types that have implicit casts. Data types are divided into several basic type categories, including boolean, numeric, string, bitstring, datetime, timespan, geometric, network, and user-defined. (For a list see Table 52.63; but note it is also possible to create custom type categories.) Within each category there can be one or more preferred types, which are preferred when there is a choice of possible types. With careful selection of preferred types and available implicit casts, it is possible to ensure that ambiguous expressions (those with multiple candidate parsing solutions) can be resolved in a useful way.

All type conversion rules are designed with several principles in mind:

- Implicit conversions should never have surprising or unpredictable outcomes.
- There should be no extra overhead in the parser or executor if a query does not need implicit type conversion. That is, if a query is well-formed and the types already match, then the query should execute without spending extra time in the parser and without introducing unnecessary implicit conversion calls in the query.
- Additionally, if a query usually requires an implicit conversion for a function, and if then the user defines a new function with the correct argument types, the parser should use this new function and no longer do implicit conversion to use the old function.

10.2. Operators

The specific operator that is referenced by an operator expression is determined using the following procedure. Note that this procedure is indirectly affected by the precedence of the operators involved, since that will determine which sub-expressions are taken to be the inputs of which operators. See Section 4.1.6 for more information.

Operator Type Resolution

1. Select the operators to be considered from the pg_operator system catalog. If a non-schema-qualified operator name was used (the usual case), the operators considered are those with the matching name and argument count that are visible in the current search path (see Section 5.8.3). If a qualified operator name was given, only operators in the specified schema are considered.

   (Optional) If the search path finds multiple operators with identical argument types, only the one appearing earliest in the path is considered. Operators with different argument types are considered on an equal footing regardless of search path position.

2. Check for an operator accepting exactly the input argument types. If one exists (there can be only one exact match in the set of operators considered), use it. Lack of an exact match creates a security hazard when calling, via qualified name (not typical), any operator found in a schema.

1 The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.
Type Conversion

that permits untrusted users to create objects. In such situations, cast arguments to force an exact match.

a. (Optional) If one argument of a binary operator invocation is of the unknown type, then assume it is the same type as the other argument for this check. Invocations involving two unknown inputs, or a unary operator with an unknown input, will never find a match at this step.

b. (Optional) If one argument of a binary operator invocation is of the unknown type and the other is of a domain type, next check to see if there is an operator accepting exactly the domain's base type on both sides; if so, use it.

3. Look for the best match.

a. Discard candidate operators for which the input types do not match and cannot be converted (using an implicit conversion) to match. unknown literals are assumed to be convertible to anything for this purpose. If only one candidate remains, use it; else continue to the next step.

b. If any input argument is of a domain type, treat it as being of the domain's base type for all subsequent steps. This ensures that domains act like their base types for purposes of ambiguous-operator resolution.

c. Run through all candidates and keep those with the most exact matches on input types. Keep all candidates if none have exact matches. If only one candidate remains, use it; else continue to the next step.

d. Run through all candidates and keep those that accept preferred types (of the input data type's type category) at the most positions where type conversion will be required. Keep all candidates if none accept preferred types. If only one candidate remains, use it; else continue to the next step.

e. If any input arguments are unknown, check the type categories accepted at those argument positions by the remaining candidates. At each position, select the string category if any candidate accepts that category. (This bias towards string is appropriate since an unknown-type literal looks like a string.) Otherwise, if all the remaining candidates accept the same type category, select that category; otherwise fail because the correct choice cannot be deduced without more clues. Now discard candidates that do not accept the selected type category. Furthermore, if any candidate accepts a preferred type in that category, discard candidates that accept non-preferred types for that argument. Keep all candidates if none survive these tests. If only one candidate remains, use it; else continue to the next step.

f. If there are both unknown and known-type arguments, and all the known-type arguments have the same type, assume that the unknown arguments are also of that type, and check which candidates can accept that type at the unknown-argument positions. If exactly one candidate passes this test, use it. Otherwise, fail.

Some examples follow.

Example 10.1. Factorial Operator Type Resolution

There is only one factorial operator (postfix !) defined in the standard catalog, and it takes an argument of type bigint. The scanner assigns an initial type of integer to the argument in this query expression:

```
SELECT 40 ! AS "40 factorial";
```

```
  40 factorial
  ------------------------
                  815915283247897734345611269596115894272000000000
```
(1 row)

So the parser does a type conversion on the operand and the query is equivalent to:

```
SELECT CAST(40 AS bigint) AS "40 factorial";
```

**Example 10.2. String Concatenation Operator Type Resolution**

A string-like syntax is used for working with string types and for working with complex extension types. Strings with unspecified type are matched with likely operator candidates.

An example with one unspecified argument:

```
SELECT text 'abc' || 'def' AS "text and unknown";
```

<table>
<thead>
<tr>
<th>text and unknown</th>
</tr>
</thead>
<tbody>
<tr>
<td>abcdef</td>
</tr>
</tbody>
</table>

(1 row)

In this case the parser looks to see if there is an operator taking `text` for both arguments. Since there is, it assumes that the second argument should be interpreted as type `text`.

Here is a concatenation of two values of unspecified types:

```
SELECT 'abc' || 'def' AS "unspecified";
```

<table>
<thead>
<tr>
<th>unspecified</th>
</tr>
</thead>
<tbody>
<tr>
<td>abcdef</td>
</tr>
</tbody>
</table>

(1 row)

In this case there is no initial hint for which type to use, since no types are specified in the query. So, the parser looks for all candidate operators and finds that there are candidates accepting both string-category and bit-string-category inputs. Since string category is preferred when available, that category is selected, and then the preferred type for strings, `text`, is used as the specific type to resolve the unknown-type literals as.

**Example 10.3. Absolute-Value and Negation Operator Type Resolution**

The PostgreSQL operator catalog has several entries for the prefix operator `@`, all of which implement absolute-value operations for various numeric data types. One of these entries is for type `float8`, which is the preferred type in the numeric category. Therefore, PostgreSQL will use that entry when faced with an `unknown` input:

```
SELECT @ '-4.5' AS "abs";
```

<table>
<thead>
<tr>
<th>abs</th>
</tr>
</thead>
<tbody>
<tr>
<td>4.5</td>
</tr>
</tbody>
</table>

(1 row)

Here the system has implicitly resolved the unknown-type literal as type `float8` before applying the chosen operator. We can verify that `float8` and not some other type was used:

```
SELECT @ '-4.5e500' AS "abs";
```

ERROR: "-4.5e500" is out of range for type double precision
On the other hand, the prefix operator \(~\) (bitwise negation) is defined only for integer data types, not for `float8`. So, if we try a similar case with \(~\), we get:

```
SELECT ~ '20' AS "negation";
```

ERROR: operator is not unique: ~ "unknown"
HINT: Could not choose a best candidate operator. You might need to add explicit type casts.

This happens because the system cannot decide which of the several possible \(~\) operators should be preferred. We can help it out with an explicit cast:

```
SELECT ~ CAST('20' AS int8) AS "negation";
```

```
<table>
<thead>
<tr>
<th>negation</th>
</tr>
</thead>
<tbody>
<tr>
<td>-21</td>
</tr>
</tbody>
</table>
```

(1 row)

**Example 10.4. Array Inclusion Operator Type Resolution**

Here is another example of resolving an operator with one known and one unknown input:

```
SELECT array[1,2] <@ '{1,2,3}' as "is subset";
```

```
<table>
<thead>
<tr>
<th>is subset</th>
</tr>
</thead>
<tbody>
<tr>
<td>t</td>
</tr>
</tbody>
</table>
```

(1 row)

The PostgreSQL operator catalog has several entries for the infix operator \(<@\), but the only two that could possibly accept an integer array on the left-hand side are array inclusion (\(\) anyarray \(<@\) anyarray) and range inclusion (\(\) anyelement \(<@\) anyrange). Since none of these polymorphic pseudo-types (see Section 8.21) are considered preferred, the parser cannot resolve the ambiguity on that basis. However, Step 3.f tells it to assume that the unknown-type literal is of the same type as the other input, that is, integer array. Now only one of the two operators can match, so array inclusion is selected. (Had range inclusion been selected, we would have gotten an error, because the string does not have the right format to be a range literal.)

**Example 10.5. Custom Operator on a Domain Type**

Users sometimes try to declare operators applying just to a domain type. This is possible but is not nearly as useful as it might seem, because the operator resolution rules are designed to select operators applying to the domain's base type. As an example consider

```
CREATE DOMAIN mytext AS text CHECK(...);
CREATE FUNCTION mytext_eq_text (mytext, text) RETURNS boolean AS ...;
CREATE OPERATOR = (procedure=mytext_eq_text, leftarg=mytext, rightarg=text);
CREATE TABLE mytable (val mytext);
SELECT * FROM mytable WHERE val = 'foo';
```

This query will not use the custom operator. The parser will first see if there is a \(\) mytext = mytext \(\) operator (Step 2.a), which there is not; then it will consider the domain's base type \(\) text, and see if
there is a `text = text` operator (Step 2.b), which there is; so it resolves the unknown-type literal as `text` and uses the `text = text` operator. The only way to get the custom operator to be used is to explicitly cast the literal:

```
SELECT * FROM mytable WHERE val = text 'foo';
```

so that the `mytext = text` operator is found immediately according to the exact-match rule. If the best-match rules are reached, they actively discriminate against operators on domain types. If they did not, such an operator would create too many ambiguous-operator failures, because the casting rules always consider a domain as castable to or from its base type, and so the domain operator would be considered usable in all the same cases as a similarly-named operator on the base type.

## 10.3. Functions

The specific function that is referenced by a function call is determined using the following procedure.

### Function Type Resolution

1. Select the functions to be considered from the `pg_proc` system catalog. If a non-schema-qualified function name was used, the functions considered are those with the matching name and argument count that are visible in the current search path (see Section 5.8.3). If a qualified function name was given, only functions in the specified schema are considered.

   a. (Optional) If the search path finds multiple functions of identical argument types, only the one appearing earliest in the path is considered. Functions of different argument types are considered on an equal footing regardless of search path position.

   b. (Optional) If a function is declared with a `VARIADIC` array parameter, and the call does not use the `VARIADIC` keyword, then the function is treated as if the array parameter were replaced by one or more occurrences of its element type, as needed to match the call. After such expansion the function might have effective argument types identical to some non-variadic function. In that case the function appearing earlier in the search path is used, or if the two functions are in the same schema, the non-variadic one is preferred.

   This creates a security hazard when calling, via qualified name\(^2\), a variadic function found in a schema that permits untrusted users to create objects. A malicious user can take control and execute arbitrary SQL functions as though you executed them. Substitute a call bearing the `VARIADIC` keyword, which bypasses this hazard. Calls populating `VARIADIC "any"` parameters often have no equivalent formulation containing the `VARIADIC` keyword. To issue those calls safely, the function's schema must permit only trusted users to create objects.

   c. (Optional) Functions that have default values for parameters are considered to match any call that omits zero or more of the defaultable parameter positions. If more than one such function matches a call, the one appearing earliest in the search path is used. If there are two or more such functions in the same schema with identical parameter types in the non-defaulted positions (which is possible if they have different sets of defaultable parameters), the system will not be able to determine which to prefer, and so an “ambiguous function call” error will result if no better match to the call can be found.

   This creates an availability hazard when calling, via qualified name\(^2\), any function found in a schema that permits untrusted users to create objects. A malicious user can create a function with the name of an existing function, replicating that function's parameters and appending novel parameters having default values. This precludes new calls to the original

\(^2\)The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.
function. To forestall this hazard, place functions in schemas that permit only trusted users to create objects.

2. Check for a function accepting exactly the input argument types. If one exists (there can be only one exact match in the set of functions considered), use it. Lack of an exact match creates a security hazard when calling, via qualified name\(^3\), a function found in a schema that permits untrusted users to create objects. In such situations, cast arguments to force an exact match. (Cases involving unknown will never find a match at this step.)

3. If no exact match is found, see if the function call appears to be a special type conversion request. This happens if the function call has just one argument and the function name is the same as the (internal) name of some data type. Furthermore, the function argument must be either an unknown-type literal, or a type that is binary-coercible to the named data type, or a type that could be converted to the named data type by applying that type's I/O functions (that is, the conversion is either to or from one of the standard string types). When these conditions are met, the function call is treated as a form of CAST specification.\(^3\)

4. Look for the best match.
   a. Discard candidate functions for which the input types do not match and cannot be converted (using an implicit conversion) to match. unknown literals are assumed to be convertible to anything for this purpose. If only one candidate remains, use it; else continue to the next step.
   b. If any input argument is of a domain type, treat it as being of the domain's base type for all subsequent steps. This ensures that domains act like their base types for purposes of ambiguous-function resolution.
   c. Run through all candidates and keep those with the most exact matches on input types. Keep all candidates if none have exact matches. If only one candidate remains, use it; else continue to the next step.
   d. Run through all candidates and keep those that accept preferred types (of the input data type's type category) at the most positions where type conversion will be required. Keep all candidates if none accept preferred types. If only one candidate remains, use it; else continue to the next step.
   e. If any input arguments are unknown, check the type categories accepted at those argument positions by the remaining candidates. At each position, select the string category if any candidate accepts that category. (This bias towards string is appropriate since an unknown-type literal looks like a string.) Otherwise, if all the remaining candidates accept the same type category, select that category; otherwise fail because the correct choice cannot be deduced without more clues. Now discard candidates that do not accept the selected type category. Furthermore, if any candidate accepts a preferred type in that category, discard candidates that accept non-preferred types for that argument. Keep all candidates if none survive these tests. If only one candidate remains, use it; else continue to the next step.
   f. If there are both unknown and known-type arguments, and all the known-type arguments have the same type, assume that the unknown arguments are also of that type, and check which candidates can accept that type at the unknown-argument positions. If exactly one candidate passes this test, use it. Otherwise, fail.

Note that the “best match” rules are identical for operator and function type resolution. Some examples follow.

\(^3\) The reason for this step is to support function-style cast specifications in cases where there is not an actual cast function. If there is a cast function, it is conventionally named after its output type, and so there is no need to have a special case. See CREATE CAST for additional commentary.
Example 10.6. Rounding Function Argument Type Resolution

There is only one `round` function that takes two arguments: it takes a first argument of type `numeric` and a second argument of type `integer`. So the following query automatically converts the first argument of type `integer` to `numeric`:

```sql
SELECT round(4, 4);
```

```
round
-------
4.0000
(1 row)
```

That query is actually transformed by the parser to:

```sql
SELECT round(CAST (4 AS numeric), 4);
```

Since numeric constants with decimal points are initially assigned the type `numeric`, the following query will require no type conversion and therefore might be slightly more efficient:

```sql
SELECT round(4.0, 4);
```

Example 10.7. Variadic Function Resolution

CREATE FUNCTION public.variadic_example(VARIADIC numeric[]) RETURNS int
   LANGUAGE sql AS 'SELECT 1';
CREATE FUNCTION

This function accepts, but does not require, the VARIADIC keyword. It tolerates both integer and numeric arguments:

```sql
SELECT public.variadic_example(0),
       public.variadic_example(0.0),
       public.variadic_example(VARIADIC array[0.0]);
```

```
variadic_example | variadic_example | variadic_example
-----------------+-----------------+------------------
       1 |                1 |                1
(1 row)
```

However, the first and second calls will prefer more-specific functions, if available:

```sql
CREATE FUNCTION public.variadic_example(numeric) RETURNS int
   LANGUAGE sql AS 'SELECT 2';
CREATE FUNCTION

CREATE FUNCTION public.variadic_example(int) RETURNS int
   LANGUAGE sql AS 'SELECT 3';
CREATE FUNCTION

SELECT public.variadic_example(0),
       public.variadic_example(0.0),
       public.variadic_example(VARIADIC array[0.0]);
```

```
variadic_example | variadic_example | variadic_example
-----------------+-----------------+------------------
       1 |                1 |                1
(1 row)
```
Given the default configuration and only the first function existing, the first and second calls are insecure. Any user could intercept them by creating the second or third function. By matching the argument type exactly and using the `VARIADIC` keyword, the third call is secure.

**Example 10.8. Substring Function Type Resolution**

There are several `substr` functions, one of which takes types `text` and `integer`. If called with a string constant of unspecified type, the system chooses the candidate function that accepts an argument of the preferred category `string` (namely of type `text`).

```sql
SELECT substr('1234', 3);
substr
--------
  34
(1 row)
```

If the string is declared to be of type `varchar`, as might be the case if it comes from a table, then the parser will try to convert it to become `text`:

```sql
SELECT substr(varchar '1234', 3);
substr
--------
  34
(1 row)
```

This is transformed by the parser to effectively become:

```sql
SELECT substr(CAST (varchar '1234' AS text), 3);
```

**Note**

The parser learns from the `pg_cast` catalog that `text` and `varchar` are binary-compatible, meaning that one can be passed to a function that accepts the other without doing any physical conversion. Therefore, no type conversion call is really inserted in this case.

And, if the function is called with an argument of type `integer`, the parser will try to convert that to `text`:

```sql
SELECT substr(1234, 3);
ERROR:  function substr(integer, integer) does not exist
HINT:  No function matches the given name and argument types. You might need to add explicit type casts.
```

This does not work because `integer` does not have an implicit cast to `text`. An explicit cast will work, however:
SELECT substr(CAST (1234 AS text), 3);

substr
--------
34
(1 row)

10.4. Value Storage

Values to be inserted into a table are converted to the destination column's data type according to the following steps.

Value Storage Type Conversion

1. Check for an exact match with the target.

2. Otherwise, try to convert the expression to the target type. This is possible if an assignment cast between the two types is registered in the pg_cast catalog (see CREATE CAST). Alternatively, if the expression is an unknown-type literal, the contents of the literal string will be fed to the input conversion routine for the target type.

3. Check to see if there is a sizing cast for the target type. A sizing cast is a cast from that type to itself. If one is found in the pg_cast catalog, apply it to the expression before storing into the destination column. The implementation function for such a cast always takes an extra parameter of type integer, which receives the destination column's atttypmod value (typically its declared length, although the interpretation of atttypmod varies for different data types), and it may take a third boolean parameter that says whether the cast is explicit or implicit. The cast function is responsible for applying any length-dependent semantics such as size checking or truncation.

Example 10.9. character Storage Type Conversion

For a target column declared as character(20) the following statement shows that the stored value is sized correctly:

```
CREATE TABLE vv (v character(20));
INSERT INTO vv SELECT 'abc' || 'def';
SELECT v, octet_length(v) FROM vv;
```

<table>
<thead>
<tr>
<th>v</th>
<th>octet_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>abcdef</td>
<td>20</td>
</tr>
</tbody>
</table>

(1 row)

What has really happened here is that the two unknown literals are resolved to text by default, allowing the || operator to be resolved as text concatenation. Then the text result of the operator is converted to bpchar (“blank-padded char”, the internal name of the character data type) to match the target column type. (Since the conversion from text to bpchar is binary-coercible, this conversion does not insert any real function call.) Finally, the sizing function bpchar(bpchar, integer, boolean) is found in the system catalog and applied to the operator's result and the stored column length. This type-specific function performs the required length check and addition of padding spaces.

10.5. UNION, CASE, and Related Constructs

SQL UNION constructs must match up possibly dissimilar types to become a single result set. The resolution algorithm is applied separately to each output column of a union query. The INTERSECT
and EXCEPT constructs resolve dissimilar types in the same way as UNION. The CASE, ARRAY, VALUES, GREATEST and LEAST constructs use the identical algorithm to match up their component expressions and select a result data type.

**Type Resolution for UNION, CASE, and Related Constructs**

1. If all inputs are of the same type, and it is not unknown, resolve as that type.
2. If any input is of a domain type, treat it as being of the domain's base type for all subsequent steps.
3. If all inputs are of type unknown, resolve as type text (the preferred type of the string category). Otherwise, unknown inputs are ignored for the purposes of the remaining rules.
4. If the non-unknown inputs are not all of the same type category, fail.
5. Choose the first non-unknown input type which is a preferred type in that category, if there is one.
6. Otherwise, choose the last non-unknown input type that allows all the preceding non-unknown inputs to be implicitly converted to it. (There always is such a type, since at least the first type in the list must satisfy this condition.)
7. Convert all inputs to the selected type. Fail if there is not a conversion from a given input to the selected type.

Some examples follow.

**Example 10.10. Type Resolution with Underspecified Types in a Union**

```sql
SELECT text 'a' AS "text" UNION SELECT 'b';
```

<table>
<thead>
<tr>
<th>text</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
</tr>
<tr>
<td>b</td>
</tr>
</tbody>
</table>

(2 rows)

Here, the unknown-type literal 'b' will be resolved to type text.

**Example 10.11. Type Resolution in a Simple Union**

```sql
SELECT 1.2 AS "numeric" UNION SELECT 1;
```

<table>
<thead>
<tr>
<th>numeric</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
</tr>
<tr>
<td>1.2</td>
</tr>
</tbody>
</table>

(2 rows)

The literal 1.2 is of type numeric, and the integer value 1 can be cast implicitly to numeric, so that type is used.

**Example 10.12. Type Resolution in a Transposed Union**

---

4 Somewhat like the treatment of domain inputs for operators and functions, this behavior allows a domain type to be preserved through a UNION or similar construct, so long as the user is careful to ensure that all inputs are implicitly or explicitly of that exact type. Otherwise the domain's base type will be preferred.
SELECT 1 AS "real" UNION SELECT CAST('2.2' AS REAL);

<table>
<thead>
<tr>
<th>real</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
</tr>
<tr>
<td>2.2</td>
</tr>
</tbody>
</table>

(2 rows)

Here, since type real cannot be implicitly cast to integer, but integer can be implicitly cast to real, the union result type is resolved as real.

**Example 10.13. Type Resolution in a Nested Union**

SELECT NULL UNION SELECT NULL UNION SELECT 1;

ERROR: UNION types text and integer cannot be matched

This failure occurs because PostgreSQL treats multiple UNIONs as a nest of pairwise operations; that is, this input is the same as

(SELECT NULL UNION SELECT NULL) UNION SELECT 1;

The inner UNION is resolved as emitting type text, according to the rules given above. Then the outer UNION has inputs of types text and integer, leading to the observed error. The problem can be fixed by ensuring that the leftmost UNION has at least one input of the desired result type.

INTERSECT and EXCEPT operations are likewise resolved pairwise. However, the other constructs described in this section consider all of their inputs in one resolution step.

### 10.6. SELECT Output Columns

The rules given in the preceding sections will result in assignment of non-unknown data types to all expressions in a SQL query, except for unspecified-type literals that appear as simple output columns of a SELECT command. For example, in

SELECT 'Hello World';

there is nothing to identify what type the string literal should be taken as. In this situation PostgreSQL will fall back to resolving the literal's type as text.

When the SELECT is one arm of a UNION (or INTERSECT or EXCEPT) construct, or when it appears within INSERT ... SELECT, this rule is not applied since rules given in preceding sections take precedence. The type of an unspecified-type literal can be taken from the other UNION arm in the first case, or from the destination column in the second case.

RETURNING lists are treated the same as SELECT output lists for this purpose.

**Note**

Prior to PostgreSQL 10, this rule did not exist, and unspecified-type literals in a SELECT output list were left as type unknown. That had assorted bad consequences, so it's been changed.
Chapter 11. Indexes

Indexes are a common way to enhance database performance. An index allows the database server to find and retrieve specific rows much faster than it could do without an index. But indexes also add overhead to the database system as a whole, so they should be used sensibly.

11.1. Introduction

Suppose we have a table similar to this:

```sql
CREATE TABLE test1 (
    id integer,
    content varchar
);
```

and the application issues many queries of the form:

```sql
SELECT content FROM test1 WHERE id = constant;
```

With no advance preparation, the system would have to scan the entire `test1` table, row by row, to find all matching entries. If there are many rows in `test1` and only a few rows (perhaps zero or one) that would be returned by such a query, this is clearly an inefficient method. But if the system has been instructed to maintain an index on the `id` column, it can use a more efficient method for locating matching rows. For instance, it might only have to walk a few levels deep into a search tree.

A similar approach is used in most non-fiction books: terms and concepts that are frequently looked up by readers are collected in an alphabetic index at the end of the book. The interested reader can scan the index relatively quickly and flip to the appropriate page(s), rather than having to read the entire book to find the material of interest. Just as it is the task of the author to anticipate the items that readers are likely to look up, it is the task of the database programmer to foresee which indexes will be useful.

The following command can be used to create an index on the `id` column, as discussed:

```sql
CREATE INDEX test1_id_index ON test1 (id);
```

The name `test1_id_index` can be chosen freely, but you should pick something that enables you to remember later what the index was for.

To remove an index, use the `DROP INDEX` command. Indexes can be added to and removed from tables at any time.

Once an index is created, no further intervention is required: the system will update the index when the table is modified, and it will use the index in queries when it thinks doing so would be more efficient than a sequential table scan. But you might have to run the `ANALYZE` command regularly to update statistics to allow the query planner to make educated decisions. See Chapter 14 for information about how to find out whether an index is used and when and why the planner might choose not to use an index.

Indexes can also benefit `UPDATE` and `DELETE` commands with search conditions. Indexes can moreover be used in join searches. Thus, an index defined on a column that is part of a join condition can also significantly speed up queries with joins.

Creating an index on a large table can take a long time. By default, PostgreSQL allows reads (`SELECT` statements) to occur on the table in parallel with index creation, but writes (`INSERT`, `UPDATE`, `DELETE`) are blocked until the index build is finished. In production environments this is often unac-
ceptable. It is possible to allow writes to occur in parallel with index creation, but there are several
caveats to be aware of — for more information see Building Indexes Concurrently.

After an index is created, the system has to keep it synchronized with the table. This adds overhead
to data manipulation operations. Therefore indexes that are seldom or never used in queries should
be removed.

11.2. Index Types

PostgreSQL provides several index types: B-tree, Hash, GiST, SP-GiST, GIN and BRIN. Each index
type uses a different algorithm that is best suited to different types of queries. By default, the CREATE INDEX command creates B-tree indexes, which fit the most common situations.

B-trees can handle equality and range queries on data that can be sorted into some ordering. In particular, the PostgreSQL query planner will consider using a B-tree index whenever an indexed column is involved in a comparison using one of these operators:

\[
< \\
\leq \\
= \\
\geq \\
> \\
\]

Constructs equivalent to combinations of these operators, such as BETWEEN and IN, can also be implemented with a B-tree index search. Also, an IS NULL or IS NOT NULL condition on an index column can be used with a B-tree index.

The optimizer can also use a B-tree index for queries involving the pattern matching operators LIKE and ~ if the pattern is a constant and is anchored to the beginning of the string — for example, \texttt{col LIKE 'foo%' or col \textasciitilde '^foo'}, but not \texttt{col LIKE '%bar'}. However, if your database does not use the C locale you will need to create the index with a special operator class to support indexing of pattern-matching queries; see Section 11.9 below. It is also possible to use B-tree indexes for \texttt{ILIKE} and \texttt{~*}, but only if the pattern starts with non-alphabetic characters, i.e., characters that are not affected by upper/lower case conversion.

B-tree indexes can also be used to retrieve data in sorted order. This is not always faster than a simple scan and sort, but it is often helpful.

Hash indexes can only handle simple equality comparisons. The query planner will consider using a hash index whenever an indexed column is involved in a comparison using the = operator. The following command is used to create a hash index:

\[
\text{CREATE INDEX name ON table USING HASH (column)}; \\
\]

GiST indexes are not a single kind of index, but rather an infrastructure within which many different indexing strategies can be implemented. Accordingly, the particular operators with which a GiST index can be used vary depending on the indexing strategy (the operator class). As an example, the standard distribution of PostgreSQL includes GiST operator classes for several two-dimensional geometric data types, which support indexed queries using these operators:

\[
<< \\
\&< \\
\&> \\
>> \\
<<| \\
\&<| \\
\&>| \\
|>> \\
\]

372
The GiST operator classes included in the standard distribution are documented in Table 64.1. Many other GiST operator classes are available in the contrib collection or as separate projects. For more information see Chapter 64.

GiST indexes are also capable of optimizing “nearest-neighbor” searches, such as

```
SELECT * FROM places ORDER BY location <-> point '(101,456)' LIMIT 10;
```

which finds the ten places closest to a given target point. The ability to do this is again dependent on the particular operator class being used. In Table 64.1, operators that can be used in this way are listed in the column “Ordering Operators”.

SP-GiST indexes, like GiST indexes, offer an infrastructure that supports various kinds of searches. SP-GiST permits implementation of a wide range of different non-balanced disk-based data structures, such as quadtrees, k-d trees, and radix trees (tries). As an example, the standard distribution of PostgreSQL includes SP-GiST operator classes for two-dimensional points, which support indexed queries using these operators:

```
<<
>>
~=
@<
^>
```

(See Section 9.11 for the meaning of these operators.) The SP-GiST operator classes included in the standard distribution are documented in Table 65.1. For more information see Chapter 65.

GIN indexes are “inverted indexes” which are appropriate for data values that contain multiple component values, such as arrays. An inverted index contains a separate entry for each component value, and can efficiently handle queries that test for the presence of specific component values.

Like GiST and SP-GiST, GIN can support many different user-defined indexing strategies, and the particular operators with which a GIN index can be used vary depending on the indexing strategy. As an example, the standard distribution of PostgreSQL includes a GIN operator class for arrays, which supports indexed queries using these operators:

```
@<
}@=
```

(See Section 9.18 for the meaning of these operators.) The GIN operator classes included in the standard distribution are documented in Table 66.1. Many other GIN operator classes are available in the contrib collection or as separate projects. For more information see Chapter 66.

BRIN indexes (a shorthand for Block Range INdexes) store summaries about the values stored in consecutive physical block ranges of a table. Like GiST, SP-GiST and GIN, BRIN can support many different indexing strategies, and the particular operators with which a BRIN index can be used vary depending on the indexing strategy. For data types that have a linear sort order, the indexed data corresponds to the minimum and maximum values of the values in the column for each block range. This supports indexed queries using these operators:
The BRIN operator classes included in the standard distribution are documented in Table 67.1. For more information see Chapter 67.

11.3. Multicolumn Indexes

An index can be defined on more than one column of a table. For example, if you have a table of this form:

```
CREATE TABLE test2 (
    major int,
    minor int,
    name varchar
);
```

(say, you keep your /dev directory in a database...) and you frequently issue queries like:

```
SELECT name FROM test2 WHERE major = constant AND minor = constant;
```

then it might be appropriate to define an index on the columns major and minor together, e.g.:

```
CREATE INDEX test2_mm_idx ON test2 (major, minor);
```

Currently, only the B-tree, GiST, GIN, and BRIN index types support multicolumn indexes. Up to 32 columns can be specified. (This limit can be altered when building PostgreSQL; see the file pg_config_manual.h.)

A multicolumn B-tree index can be used with query conditions that involve any subset of the index's columns, but the index is most efficient when there are constraints on the leading (leftmost) columns. The exact rule is that equality constraints on leading columns, plus any inequality constraints on the first column that does not have an equality constraint, will be used to limit the portion of the index that is scanned. Constraints on columns to the right of these columns are checked in the index, so they save visits to the table proper, but they do not reduce the portion of the index that has to be scanned. For example, given an index on (a, b, c) and a query condition WHERE a = 5 AND b >= 42 AND c < 77, the index would have to be scanned from the first entry with a = 5 and b = 42 up through the last entry with a = 5. Index entries with c >= 77 would be skipped, but they'd still have to be scanned through. This index could in principle be used for queries that have constraints on b and/or c with no constraint on a — but the entire index would have to be scanned, so in most cases the planner would prefer a sequential table scan over using the index.

A multicolumn GiST index can be used with query conditions that involve any subset of the index's columns. Conditions on additional columns restrict the entries returned by the index, but the condition on the first column is the most important one for determining how much of the index needs to be scanned. A GiST index will be relatively ineffective if its first column has only a few distinct values, even if there are many distinct values in additional columns.

A multicolumn GIN index can be used with query conditions that involve any subset of the index's columns. Unlike B-tree or GiST, index search effectiveness is the same regardless of which index column(s) the query conditions use.

A multicolumn BRIN index can be used with query conditions that involve any subset of the index's columns. Like GIN and unlike B-tree or GiST, index search effectiveness is the same regardless of
which index column(s) the query conditions use. The only reason to have multiple BRIN indexes instead of one multicol

Of course, each column must be used with operators appropriate to the index type; clauses that involve other operators will not be considered.

Multicolumn indexes should be used sparingly. In most situations, an index on a single column is sufficient and saves space and time. Indexes with more than three columns are unlikely to be helpful unless the usage of the table is extremely stylized. See also Section 11.5 and Section 11.11 for some discussion of the merits of different index configurations.

11.4. Indexes and ORDER BY

In addition to simply finding the rows to be returned by a query, an index may be able to deliver them in a specific sorted order. This allows a query’s ORDER BY specification to be honored without a separate sorting step. Of the index types currently supported by PostgreSQL, only B-tree can produce sorted output — the other index types return matching rows in an unspecified, implementation-dependent order.

The planner will consider satisfying an ORDER BY specification either by scanning an available index that matches the specification, or by scanning the table in physical order and doing an explicit sort. For a query that requires scanning a large fraction of the table, an explicit sort is likely to be faster than using an index because it requires less disk I/O due to following a sequential access pattern. Indexes are more useful when only a few rows need be fetched. An important special case is ORDER BY in combination with LIMIT n: an explicit sort will have to process all the data to identify the first n rows, but if there is an index matching the ORDER BY, the first n rows can be retrieved directly, without scanning the remainder at all.

By default, B-tree indexes store their entries in ascending order with nulls last. This means that a forward scan of an index on column x produces output satisfying ORDER BY x (or more verbosely, ORDER BY x ASC NULLS LAST). The index can also be scanned backward, producing output satisfying ORDER BY x DESC (or more verbosely, ORDER BY x DESC NULLS FIRST, since NULLS FIRST is the default for ORDER BY DESC).

You can adjust the ordering of a B-tree index by including the options ASC, DESC, NULLS FIRST, and/or NULLS LAST when creating the index; for example:

CREATE INDEX test2_info_nulls_low ON test2 (info NULLS FIRST);
CREATE INDEX test3_desc_index ON test3 (id DESC NULLS LAST);

An index stored in ascending order with nulls first can satisfy either ORDER BY x ASC NULLS FIRST or ORDER BY x DESC NULLS LAST depending on which direction it is scanned in.

You might wonder why bother providing all four options, when two options together with the possibility of backward scan would cover all the variants of ORDER BY. In single-column indexes the options are indeed redundant, but in multicol

Obviously, indexes with non-default sort orderings are a fairly specialized feature, but sometimes they can produce tremendous speedups for certain queries. Whether it’s worth maintaining such an index depends on how often you use queries that require a special sort ordering.

11.5. Combining Multiple Indexes
Indexes

A single index scan can only use query clauses that use the index's columns with operators of its operator class and are joined with AND. For example, given an index on (a, b) a query condition like WHERE \( a = 5 \) AND \( b = 6 \) could use the index, but a query like WHERE \( a = 5 \) OR \( b = 6 \) could not directly use the index.

Fortunately, PostgreSQL has the ability to combine multiple indexes (including multiple uses of the same index) to handle cases that cannot be implemented by single index scans. The system can form AND and OR conditions across several index scans. For example, a query like WHERE \( x = 42 \) OR \( x = 47 \) OR \( x = 53 \) OR \( x = 99 \) could be broken down into four separate scans of an index on \( x \), each scan using one of the query clauses. The results of these scans are then ORed together to produce the result. Another example is that if we have separate indexes on \( x \) and \( y \), one possible implementation of a query like WHERE \( x = 5 \) AND \( y = 6 \) is to use each index with the appropriate query clause and then AND together the index results to identify the result rows.

To combine multiple indexes, the system scans each needed index and prepares a bitmap in memory giving the locations of table rows that are reported as matching that index's conditions. The bitmaps are then ANDed and ORed together as needed by the query. Finally, the actual table rows are visited and returned. The table rows are visited in physical order, because that is how the bitmap is laid out; this means that any ordering of the original indexes is lost, and so a separate sort step will be needed if the query has an ORDER BY clause. For this reason, and because each additional index scan adds extra time, the planner will sometimes choose to use a simple index scan even though additional indexes are available that could have been used as well.

In all but the simplest applications, there are various combinations of indexes that might be useful, and the database developer must make trade-offs to decide which indexes to provide. Sometimes multicolumn indexes are best, but sometimes it's better to create separate indexes and rely on the index-combination feature. For example, if your workload includes a mix of queries that sometimes involve only column \( x \), sometimes only column \( y \), and sometimes both columns, you might choose to create two separate indexes on \( x \) and \( y \), relying on index combination to process the queries that use both columns. You could also create a multicolumn index on \((x, y)\). This index would typically be more efficient than index combination for queries involving both columns, but as discussed in Section 11.3, it would be almost useless for queries involving only \( y \), so it should not be the only index. A combination of the multicolumn index and a separate index on \( y \) would serve reasonably well. For queries involving only \( x \), the multicolumn index could be used, though it would be larger and hence slower than an index on \( x \) alone. The last alternative is to create all three indexes, but this is probably only reasonable if the table is searched much more often than it is updated and all three types of query are common. If one of the types of query is much less common than the others, you'd probably settle for creating just the two indexes that best match the common types.

11.6. Unique Indexes

Indexes can also be used to enforce uniqueness of a column's value, or the uniqueness of the combined values of more than one column.

```
CREATE UNIQUE INDEX name ON table (column [, ...])
[ INCLUDE (column [, ...]) ];
```

Currently, only B-tree indexes can be declared unique.

When an index is declared unique, multiple table rows with equal indexed values are not allowed. Null values are not considered equal. A multicolumn unique index will only reject cases where all indexed columns are equal. Columns listed in the INCLUDE clause, if any, aren't considered when determining whether index entries are equal.

PostgreSQL automatically creates a unique index when a unique constraint or primary key is defined for a table. The index covers the columns that make up the primary key or unique constraint (a multicolumn index, if appropriate), and is the mechanism that enforces the constraint.
Indexes

11.7. Indexes on Expressions

An index column need not be just a column of the underlying table, but can be a function or scalar expression computed from one or more columns of the table. This feature is useful to obtain fast access to tables based on the results of computations.

For example, a common way to do case-insensitive comparisons is to use the `lower` function:

```
SELECT * FROM test1 WHERE lower(col1) = 'value';
```

This query can use an index if one has been defined on the result of the `lower(col1)` function:

```
CREATE INDEX test1_lower_col1_idx ON test1 (lower(col1));
```

If we were to declare this index `UNIQUE`, it would prevent creation of rows whose `col1` values differ only in case, as well as rows whose `col1` values are actually identical. Thus, indexes on expressions can be used to enforce constraints that are not definable as simple unique constraints.

As another example, if one often does queries like:

```
SELECT * FROM people WHERE (first_name || ' ' || last_name) = 'John Smith';
```

then it might be worth creating an index like this:

```
CREATE INDEX people_names ON people ((first_name || ' ' || last_name));
```

The syntax of the `CREATE INDEX` command normally requires writing parentheses around index expressions, as shown in the second example. The parentheses can be omitted when the expression is just a function call, as in the first example.

Index expressions are relatively expensive to maintain, because the derived expression(s) must be computed for each row upon insertion and whenever it is updated. However, the index expressions are not recomputed during an indexed search, since they are already stored in the index. In both examples above, the system sees the query as just `WHERE indexedcolumn = 'constant'` and so the speed of the search is equivalent to any other simple index query. Thus, indexes on expressions are useful when retrieval speed is more important than insertion and update speed.

11.8. Partial Indexes

A partial index is an index built over a subset of a table; the subset is defined by a conditional expression (called the predicate of the partial index). The index contains entries only for those table rows that satisfy the predicate. Partial indexes are a specialized feature, but there are several situations in which they are useful.

One major reason for using a partial index is to avoid indexing common values. Since a query searching for a common value (one that accounts for more than a few percent of all the table rows) will not use
the index anyway, there is no point in keeping those rows in the index at all. This reduces the size of the index, which will speed up those queries that do use the index. It will also speed up many table update operations because the index does not need to be updated in all cases. Example 11.1 shows a possible application of this idea.

Example 11.1. Setting up a Partial Index to Exclude Common Values

Suppose you are storing web server access logs in a database. Most accesses originate from the IP address range of your organization but some are from elsewhere (say, employees on dial-up connections). If your searches by IP are primarily for outside accesses, you probably do not need to index the IP range that corresponds to your organization’s subnet.

Assume a table like this:

```sql
CREATE TABLE access_log (
    url varchar,
    client_ip inet,
    ...
);
```

To create a partial index that suits our example, use a command such as this:

```sql
CREATE INDEX access_log_client_ip_ix ON access_log (client_ip)
WHERE NOT (client_ip > inet '192.168.100.0' AND client_ip < inet '192.168.100.255');
```

A typical query that can use this index would be:

```sql
SELECT *
FROM access_log
WHERE url = '/index.html' AND client_ip = inet '212.78.10.32';
```

A query that cannot use this index is:

```sql
SELECT *
FROM access_log
WHERE client_ip = inet '192.168.100.23';
```

Observe that this kind of partial index requires that the common values be predetermined, so such partial indexes are best used for data distributions that do not change. The indexes can be recreated occasionally to adjust for new data distributions, but this adds maintenance effort.

Another possible use for a partial index is to exclude values from the index that the typical query workload is not interested in; this is shown in Example 11.2. This results in the same advantages as listed above, but it prevents the “uninteresting” values from being accessed via that index, even if an index scan might be profitable in that case. Obviously, setting up partial indexes for this kind of scenario will require a lot of care and experimentation.

Example 11.2. Setting up a Partial Index to Exclude Uninteresting Values

If you have a table that contains both billed and unbilled orders, where the unbilled orders take up a small fraction of the total table and yet those are the most-accessed rows, you can improve performance by creating an index on just the unbilled rows. The command to create the index would look like this:
Indexes

CREATE INDEX orders_unbilled_index ON orders (order_nr)
    WHERE billed is not true;

A possible query to use this index would be:

SELECT * FROM orders WHERE billed is not true AND order_nr < 10000;

However, the index can also be used in queries that do not involve order_nr at all, e.g.:

SELECT * FROM orders WHERE billed is not true AND amount > 5000.00;

This is not as efficient as a partial index on the amount column would be, since the system has to scan the entire index. Yet, if there are relatively few unbilled orders, using this partial index just to find the unbilled orders could be a win.

Note that this query cannot use this index:

SELECT * FROM orders WHERE order_nr = 3501;

The order 3501 might be among the billed or unbilled orders.

Example 11.2 also illustrates that the indexed column and the column used in the predicate do not need to match. PostgreSQL supports partial indexes with arbitrary predicates, so long as only columns of the table being indexed are involved. However, keep in mind that the predicate must match the conditions used in the queries that are supposed to benefit from the index. To be precise, a partial index can be used in a query only if the system can recognize that the WHERE condition of the query mathematically implies the predicate of the index. PostgreSQL does not have a sophisticated theorem prover that can recognize mathematically equivalent expressions that are written in different forms. (Not only is such a general theorem prover extremely difficult to create, it would probably be too slow to be of any real use.) The system can recognize simple inequality implications, for example “x < 1” implies “x < 2”; otherwise the predicate condition must exactly match part of the query’s WHERE condition or the index will not be recognized as usable. Matching takes place at query planning time, not at run time. As a result, parameterized query clauses do not work with a partial index. For example a prepared query with a parameter might specify “x < ?” which will never imply “x < 2” for all possible values of the parameter.

A third possible use for partial indexes does not require the index to be used in queries at all. The idea here is to create a unique index over a subset of a table, as in Example 11.3. This enforces uniqueness among the rows that satisfy the index predicate, without constraining those that do not.

Example 11.3. Setting up a Partial Unique Index

Suppose that we have a table describing test outcomes. We wish to ensure that there is only one “successful” entry for a given subject and target combination, but there might be any number of “unsuccesful” entries. Here is one way to do it:

CREATE TABLE tests (  
    subject text,  
    target text,  
    success boolean,  
    ...  
);  

CREATE UNIQUE INDEX tests_success_constraint ON tests (subject, target)  
    WHERE success;
This is a particularly efficient approach when there are few successful tests and many unsuccessful ones.

Finally, a partial index can also be used to override the system's query plan choices. Also, data sets with peculiar distributions might cause the system to use an index when it really should not. In that case the index can be set up so that it is not available for the offending query. Normally, PostgreSQL makes reasonable choices about index usage (e.g., it avoids them when retrieving common values, so the earlier example really only saves index size, it is not required to avoid index usage), and grossly incorrect plan choices are cause for a bug report.

Keep in mind that setting up a partial index indicates that you know at least as much as the query planner knows, in particular you know when an index might be profitable. Forming this knowledge requires experience and understanding of how indexes in PostgreSQL work. In most cases, the advantage of a partial index over a regular index will be minimal.

More information about partial indexes can be found in [ston89b], [olson93], and [seshadri95].

## 11.9. Operator Classes and Operator Families

An index definition can specify an operator class for each column of an index.

```
CREATE INDEX name ON table (column opclass [sort options] [, ...]);
```

The operator class identifies the operators to be used by the index for that column. For example, a B-tree index on the type `int4` would use the `int4_ops` class; this operator class includes comparison functions for values of type `int4`. In practice the default operator class for the column's data type is usually sufficient. The main reason for having operator classes is that for some data types, there could be more than one meaningful index behavior. For example, we might want to sort a complex-number data type either by absolute value or by real part. We could do this by defining two operator classes for the data type and then selecting the proper class when making an index. The operator class determines the basic sort ordering (which can then be modified by adding sort options `COLLATE`, `ASC/DESC` and/or `NULLS FIRST/NULLS LAST`).

There are also some built-in operator classes besides the default ones:

- The operator classes `text_pattern_ops`, `varchar_pattern_ops`, and `bpchar_pattern_ops` support B-tree indexes on the types `text`, `varchar`, and `char` respectively. The difference from the default operator classes is that the values are compared strictly character by character rather than according to the locale-specific collation rules. This makes these operator classes suitable for use by queries involving pattern matching expressions (LIKE or POSIX regular expressions) when the database does not use the standard “C” locale. As an example, you might index a `varchar` column like this:

  ```
  CREATE INDEX test_index ON test_table (col varchar_pattern_ops);
  ```

  Note that you should also create an index with the default operator class if you want queries involving ordinary `<`, `<=`, `>`, or `>=` comparisons to use an index. Such queries cannot use the `xxx_pattern_ops` operator classes. (Ordinary equality comparisons can use these operator classes, however.) It is possible to create multiple indexes on the same column with different operator classes. If you do use the C locale, you do not need the `xxx_pattern_ops` operator classes, because an index with the default operator class is usable for pattern-matching queries in the C locale.

The following query shows all defined operator classes:

```
SELECT am.amname AS index_method,
       opc.opcname AS opclass_name,
```
An operator class is actually just a subset of a larger structure called an operator family. In cases where several data types have similar behaviors, it is frequently useful to define cross-data-type operators and allow these to work with indexes. To do this, the operator classes for each of the types must be grouped into the same operator family. The cross-type operators are members of the family, but are not associated with any single class within the family.

This expanded version of the previous query shows the operator family each operator class belongs to:

```
SELECT am.amname AS index_method,
       opc.opcname AS opclass_name,
       opf.opfname AS opfamily_name,
       opc.opcintype::regtype AS indexed_type,
       opc.opcdefault AS is_default
FROM pg_am am, pg_opclass opc, pg_opfamily opf
WHERE opc.opcmethod = am.oid AND
       opc.opcfamily = opf.oid
ORDER BY index_method, opclass_name;
```

This query shows all defined operator families and all the operators included in each family:

```
SELECT am.amname AS index_method,
       opf.opfname AS opfamily_name,
       amop.amopopr::regoperator AS opfamily_operator
FROM pg_am am, pg_opfamily opf, pg_amop amop
WHERE opf.opfmethod = am.oid AND
       amop.amopfamily = opf.oid
ORDER BY index_method, opfamily_name, opfamily_operator;
```

### 11.10. Indexes and Collations

An index can support only one collation per index column. If multiple collations are of interest, multiple indexes may be needed.

Consider these statements:

```
CREATE TABLE test1c (  
id integer,  
    content varchar COLLATE "x"
);  
CREATE INDEX test1c_content_index ON test1c (content);  
```

The index automatically uses the collation of the underlying column. So a query of the form

```
SELECT * FROM test1c WHERE content > constant;
```

could use the index, because the comparison will by default use the collation of the column. However, this index cannot accelerate queries that involve some other collation. So if queries of the form, say,
Indexes


df5 Selector * FROM test1c WHERE content > constant COLLATE "y";

are also of interest, an additional index could be created that supports the "y" collation, like this:

CREATE INDEX test1c_content_y_index ON test1c (content COLLATE "y");

11.1. Index-Only Scans

All indexes in PostgreSQL are secondary indexes, meaning that each index is stored separately from the table's main data area (which is called the table's heap in PostgreSQL terminology). This means that in an ordinary index scan, each row retrieval requires fetching data from both the index and the heap. Furthermore, while the index entries that match a given indexable WHERE condition are usually close together in the index, the table rows they reference might be anywhere in the heap. The heap-access portion of an index scan thus involves a lot of random access into the heap, which can be slow, particularly on traditional rotating media. (As described in Section 11.5, bitmap scans try to alleviate this cost by doing the heap accesses in sorted order, but that only goes so far.)

To solve this performance problem, PostgreSQL supports index-only scans, which can answer queries from an index alone without any heap access. The basic idea is to return values directly out of each index entry instead of consulting the associated heap entry. There are two fundamental restrictions on when this method can be used:

1. The index type must support index-only scans. B-tree indexes always do. GiST and SP-GiST indexes support index-only scans for some operator classes but not others. Other index types have no support. The underlying requirement is that the index must physically store, or else be able to reconstruct, the original data value for each index entry. As a counterexample, GIN indexes cannot support index-only scans because each index entry typically holds only part of the original data value.

2. The query must reference only columns stored in the index. For example, given an index on columns x and y of a table that also has a column z, these queries could use index-only scans:

```
SELECT x, y FROM tab WHERE x = 'key';
SELECT x FROM tab WHERE x = 'key' AND y < 42;
```

but these queries could not:

```
SELECT x, z FROM tab WHERE x = 'key';
SELECT x FROM tab WHERE x = 'key' AND z < 42;
```

(Expression indexes and partial indexes complicate this rule, as discussed below.)

If these two fundamental requirements are met, then all the data values required by the query are available from the index, so an index-only scan is physically possible. But there is an additional requirement for any table scan in PostgreSQL: it must verify that each retrieved row be “visible” to the query’s MVCC snapshot, as discussed in Chapter 13. Visibility information is not stored in index entries, only in heap entries; so at first glance it would seem that every row retrieval would require a heap access anyway. And this is indeed the case, if the table row has been modified recently. However, for seldom-changing data there is a way around this problem. PostgreSQL tracks, for each page in a table’s heap, whether all rows stored in that page are old enough to be visible to all current and future transactions. This information is stored in a bit in the table’s visibility map. An index-only scan, after finding a candidate index entry, checks the visibility map bit for the corresponding heap page. If it’s set, the row is known visible and so the data can be returned with no further work. If it’s not set, the heap entry must be visited to find out whether it’s visible, so no performance advantage is gained over a standard index scan. Even in the successful case, this approach trades visibility map accesses...
Indexes

for heap accesses; but since the visibility map is four orders of magnitude smaller than the heap it
describes, far less physical I/O is needed to access it. In most situations the visibility map remains
cached in memory all the time.

In short, while an index-only scan is possible given the two fundamental requirements, it will be a
win only if a significant fraction of the table's heap pages have their all-visible map bits set. But tables
in which a large fraction of the rows are unchanging are common enough to make this type of scan
very useful in practice.

To make effective use of the index-only scan feature, you might choose to create indexes in which only
the leading columns are meant to match WHERE clauses, while the trailing columns hold “payload”
data to be returned by a query. For example, if you commonly run queries like

```
SELECT y FROM tab WHERE x = 'key';
```

the traditional approach to speeding up such queries would be to create an index on x only. However,
an index on (x, y) would offer the possibility of implementing this query as an index-only scan.
As previously discussed, such an index would be larger and hence more expensive than an index on
x alone, so this is attractive only if the table is known to be mostly static. Note it's important that the
index be declared on (x, y) not (y, x), as for most index types (particularly B-trees) searches
that do not constrain the leading index columns are not very efficient.

In principle, index-only scans can be used with expression indexes. For example, given an index on
f(x) where x is a table column, it should be possible to execute

```
SELECT f(x) FROM tab WHERE f(x) < 1;
```

as an index-only scan; and this is very attractive if f() is an expensive-to-compute function. However,
PostgreSQL’s planner is currently not very smart about such cases. It considers a query to be potentially
executable by index-only scan only when all columns needed by the query are available from the index.
In this example, x is not needed except in the context f(x), but the planner does not notice that and
concludes that an index-only scan is not possible. If an index-only scan seems sufficiently worthwhile,
this can be worked around by declaring the index to be on (f(x), x), where the second column
is not expected to be used in practice but is just there to convince the planner that an index-only scan
is possible. An additional caveat, if the goal is to avoid recalculating f(x), is that the planner won’t
necessarily match uses of f(x) that aren’t in indexable WHERE clauses to the index column. It will
usually get this right in simple queries such as shown above, but not in queries that involve joins.
These deficiencies may be remedied in future versions of PostgreSQL.

Partial indexes also have interesting interactions with index-only scans. Consider the partial index
shown in Example 11.3:

```
CREATE UNIQUE INDEX tests_success_constraint ON tests (subject,
        target)
        WHERE success;
```

In principle, we could do an index-only scan on this index to satisfy a query like

```
SELECT target FROM tests WHERE subject = 'some-subject' AND
        success;
```

But there's a problem: the WHERE clause refers to success which is not available as a result column
of the index. Nonetheless, an index-only scan is possible because the plan does not need to recheck
that part of the WHERE clause at run time: all entries found in the index necessarily have success
= true so this need not be explicitly checked in the plan. PostgreSQL versions 9.6 and later will
recognize such cases and allow index-only scans to be generated, but older versions will not.
11.12. Examining Index Usage

Although indexes in PostgreSQL do not need maintenance or tuning, it is still important to check which indexes are actually used by the real-life query workload. Examining index usage for an individual query is done with the EXPLAIN command; its application for this purpose is illustrated in Section 14.1. It is also possible to gather overall statistics about index usage in a running server, as described in Section 28.2.

It is difficult to formulate a general procedure for determining which indexes to create. There are a number of typical cases that have been shown in the examples throughout the previous sections. A good deal of experimentation is often necessary. The rest of this section gives some tips for that:

- Always run ANALYZE first. This command collects statistics about the distribution of the values in the table. This information is required to estimate the number of rows returned by a query, which is needed by the planner to assign realistic costs to each possible query plan. In absence of any real statistics, some default values are assumed, which are almost certain to be inaccurate. Examining an application's index usage without having run ANALYZE is therefore a lost cause. See Section 24.1.3 and Section 24.1.6 for more information.

- Use real data for experimentation. Using test data for setting up indexes will tell you what indexes you need for the test data, but that is all.

It is especially fatal to use very small test data sets. While selecting 1000 out of 100000 rows could be a candidate for an index, selecting 1 out of 100 rows will hardly be, because the 100 rows probably fit within a single disk page, and there is no plan that can beat sequentially fetching 1 disk page.

Also be careful when making up test data, which is often unavoidable when the application is not yet in production. Values that are very similar, completely random, or inserted in sorted order will skew the statistics away from the distribution that real data would have.

- When indexes are not used, it can be useful for testing to force their use. There are run-time parameters that can turn off various plan types (see Section 19.7.1). For instance, turning off sequential scans (enable_seqscan) and nested-loop joins (enable_nestloop), which are the most basic plans, will force the system to use a different plan. If the system still chooses a sequential scan or nested-loop join then there is probably a more fundamental reason why the index is not being used; for example, the query condition does not match the index. (What kind of query can use what kind of index is explained in the previous sections.)

- If forcing index usage does use the index, then there are two possibilities: Either the system is right and using the index is indeed not appropriate, or the cost estimates of the query plans are not reflecting reality. So you should time your query with and without indexes. The EXPLAIN ANALYZE command can be useful here.

- If it turns out that the cost estimates are wrong, there are, again, two possibilities. The total cost is computed from the per-row costs of each plan node times the selectivity estimate of the plan node. The costs estimated for the plan nodes can be adjusted via run-time parameters (described in Section 19.7.2). An inaccurate selectivity estimate is due to insufficient statistics. It might be possible to improve this by tuning the statistics-gathering parameters (see ALTER TABLE).

If you do not succeed in adjusting the costs to be more appropriate, then you might have to resort to forcing index usage explicitly. You might also want to contact the PostgreSQL developers to examine the issue.
Chapter 12. Full Text Search

12.1. Introduction

Full Text Searching (or just text search) provides the capability to identify natural-language documents that satisfy a query, and optionally to sort them by relevance to the query. The most common type of search is to find all documents containing given query terms and return them in order of their similarity to the query. Notions of query and similarity are very flexible and depend on the specific application. The simplest search considers query as a set of words and similarity as the frequency of query words in the document.

Textual search operators have existed in databases for years. PostgreSQL has ~, ~*, LIKE, and ILIKE operators for textual data types, but they lack many essential properties required by modern information systems:

• There is no linguistic support, even for English. Regular expressions are not sufficient because they cannot easily handle derived words, e.g., satisfies and satisfy. You might miss documents that contain satisfies, although you probably would like to find them when searching for satisfy. It is possible to use OR to search for multiple derived forms, but this is tedious and error-prone (some words can have several thousand derivatives).
• They provide no ordering (ranking) of search results, which makes them ineffective when thousands of matching documents are found.
• They tend to be slow because there is no index support, so they must process all documents for every search.

Full text indexing allows documents to be preprocessed and an index saved for later rapid searching. Preprocessing includes:

*Parsing documents into tokens.* It is useful to identify various classes of tokens, e.g., numbers, words, complex words, email addresses, so that they can be processed differently. In principle token classes depend on the specific application, but for most purposes it is adequate to use a predefined set of classes. PostgreSQL uses a parser to perform this step. A standard parser is provided, and custom parsers can be created for specific needs.

*Converting tokens into lexemes.* A lexeme is a string, just like a token, but it has been normalized so that different forms of the same word are made alike. For example, normalization almost always includes folding upper-case letters to lower-case, and often involves removal of suffixes (such as s or es in English). This allows searches to find variant forms of the same word, without tediously entering all the possible variants. Also, this step typically eliminates stop words, which are words that are so common that they are useless for searching. (In short, then, tokens are raw fragments of the document text, while lexemes are words that are believed useful for indexing and searching.) PostgreSQL uses dictionaries to perform this step. Various standard dictionaries are provided, and custom ones can be created for specific needs.

*Storing preprocessed documents optimized for searching.* For example, each document can be represented as a sorted array of normalized lexemes. Along with the lexemes it is often desirable to store positional information to use for proximity ranking, so that a document that contains a more “dense” region of query words is assigned a higher rank than one with scattered query words.

Dictionaries allow fine-grained control over how tokens are normalized. With appropriate dictionaries, you can:

• Define stop words that should not be indexed.
• Map synonyms to a single word using Ispell.
• Map phrases to a single word using a thesaurus.
• Map different variations of a word to a canonical form using an Ispell dictionary.
• Map different variations of a word to a canonical form using Snowball stemmer rules.

A data type tsvector is provided for storing preprocessed documents, along with a type tsquery for representing processed queries (Section 8.11). There are many functions and operators available for these data types (Section 9.13), the most important of which is the match operator @@, which we introduce in Section 12.1.2. Full text searches can be accelerated using indexes (Section 12.9).

12.1.1. What Is a Document?

A document is the unit of searching in a full text search system; for example, a magazine article or email message. The text search engine must be able to parse documents and store associations of lexemes (key words) with their parent document. Later, these associations are used to search for documents that contain query words.

For searches within PostgreSQL, a document is normally a textual field within a row of a database table, or possibly a combination (concatenation) of such fields, perhaps stored in several tables or obtained dynamically. In other words, a document can be constructed from different parts for indexing and it might not be stored anywhere as a whole. For example:

```
SELECT title || ' ' || author || ' ' || abstract || ' ' || body
AS document
FROM messages
WHERE mid = 12;

SELECT m.title || ' ' || m.author || ' ' || m.abstract || ' ' || d.body AS document
FROM messages m, docs d
WHERE mid = did AND mid = 12;
```

Note

Actually, in these example queries, coalesce should be used to prevent a single NULL attribute from causing a NULL result for the whole document.

Another possibility is to store the documents as simple text files in the file system. In this case, the database can be used to store the full text index and to execute searches, and some unique identifier can be used to retrieve the document from the file system. However, retrieving files from outside the database requires superuser permissions or special function support, so this is usually less convenient than keeping all the data inside PostgreSQL. Also, keeping everything inside the database allows easy access to document metadata to assist in indexing and display.

For text search purposes, each document must be reduced to the preprocessed tsvector format. Searching and ranking are performed entirely on the tsvector representation of a document — the original text need only be retrieved when the document has been selected for display to a user. We therefore often speak of the tsvector as being the document, but of course it is only a compact representation of the full document.

12.1.2. Basic Text Matching

Full text searching in PostgreSQL is based on the match operator @@, which returns true if the tsvector (document) matches a tsquery (query). It doesn't matter which data type is written first:

```
SELECT 'a fat cat sat on a mat and ate a fat rat':tsvector @@ 'cat & rat':tsquery;
```
As the above example suggests, a tsquery is not just raw text, any more than a tsvector is. A tsquery contains search terms, which must be already-normalized lexemes, and may combine multiple terms using AND, OR, NOT, and FOLLOWED BY operators. (For syntax details see Section 8.11.2.) There are functions to_tsquery, plainto_tsquery, and phraseto_tsquery that are helpful in converting user-written text into a proper tsquery, primarily by normalizing words appearing in the text. Similarly, to_tsvector is used to parse and normalize a document string. So in practice a text search match would look more like this:

```
SELECT to_tsvector('fat cats ate fat rats') @@ to_tsquery('fat & rat');
```

Observe that this match would not succeed if written as

```
SELECT 'fat cats ate fat rats'::tsvector @@ to_tsquery('fat & rat');
```

since here no normalization of the word rats will occur. The elements of a tsvector are lexemes, which are assumed already normalized, so rats does not match rat.

The @@ operator also supports text input, allowing explicit conversion of a text string to tsvector or tsquery to be skipped in simple cases. The variants available are:

- tsvector @@ tsquery
- tsquery @@ tsvector
- text @@ tsquery
- text @@ text

The first two of these we saw already. The form text @@ tsquery is equivalent to to_tsvector(x) @@ y. The form text @@ text is equivalent to to_tsvector(x) @@ plainto_tsquery(y).

Within a tsquery, the & (AND) operator specifies that both its arguments must appear in the document to have a match. Similarly, the | (OR) operator specifies that at least one of its arguments must appear, while the ! (NOT) operator specifies that its argument must not appear in order to have a match. For example, the query fat & ! rat matches documents that contain fat but not rat.

Searching for phrases is possible with the help of the <-> (FOLLOWED BY) tsquery operator, which matches only if its arguments have matches that are adjacent and in the given order. For example:

```
SELECT to_tsvector('fatal error') @@ to_tsquery('fatal <-> error');
```
There is a more general version of the FOLLOWED BY operator having the form <N>, where N is an integer standing for the difference between the positions of the matching lexemes. <1> is the same as <->, while <2> allows exactly one other lexeme to appear between the matches, and so on. The phraseto_tsquery function makes use of this operator to construct a tsquery that can match a multi-word phrase when some of the words are stop words. For example:

```
SELECT phraseto_tsquery('cats ate rats');
phraseto_tsquery
-----------------------------------
'cat' <-> 'ate' <-> 'rat'
```

```
SELECT phraseto_tsquery('the cats ate the rats');
phraseto_tsquery
-----------------------------------
'cat' <-> 'ate' <2> 'rat'
```

A special case that's sometimes useful is that <0> can be used to require that two patterns match the same word.

Parentheses can be used to control nesting of the tsquery operators. Without parentheses, | binds least tightly, then &, then <->, and ! most tightly.

It's worth noticing that the AND/OR/NOT operators mean something subtly different when they are within the arguments of a FOLLOWED BY operator than when they are not, because within FOLLOWED BY the exact position of the match is significant. For example, normally !x matches only documents that do not contain x anywhere. But !x <-> y matches y if it is not immediately after an x; an occurrence of x elsewhere in the document does not prevent a match. Another example is that x & y normally only requires that x and y both appear somewhere in the document, but (x & y) <-> z requires x and y to match at the same place, immediately before a z. Thus this query behaves differently from x <-> z & y <-> z, which will match a document containing two separate sequences x z and y z. (This specific query is useless as written, since x and y could not match at the same place; but with more complex situations such as prefix-match patterns, a query of this form could be useful.)

## 12.1.3. Configurations

The above are all simple text search examples. As mentioned before, full text search functionality includes the ability to do many more things: skip indexing certain words (stop words), process synonyms, and use sophisticated parsing, e.g., parse based on more than just white space. This functionality is controlled by text search configurations. PostgreSQL comes with predefined configurations for many languages, and you can easily create your own configurations. (psql's \dF command shows all available configurations.)

During installation an appropriate configuration is selected and default_text_search_config is set accordingly in postgresql.conf. If you are using the same text search configuration for the entire cluster you can use the value in postgresql.conf. To use different configurations throughout the cluster but the same configuration within any one database, use ALTER DATABASE ... SET. Otherwise, you can set default_text_search_config in each session.
Each text search function that depends on a configuration has an optional `regconfig` argument, so that the configuration to use can be specified explicitly. `default_text_search_config` is used only when this argument is omitted.

To make it easier to build custom text search configurations, a configuration is built up from simpler database objects. PostgreSQL’s text search facility provides four types of configuration-related database objects:

- **Text search parsers** break documents into tokens and classify each token (for example, as words or numbers).
- **Text search dictionaries** convert tokens to normalized form and reject stop words.
- **Text search templates** provide the functions underlying dictionaries. (A dictionary simply specifies a template and a set of parameters for the template.)
- **Text search configurations** select a parser and a set of dictionaries to use to normalize the tokens produced by the parser.

Text search parsers and templates are built from low-level C functions; therefore it requires C programming ability to develop new ones, and superuser privileges to install one into a database. (There are examples of add-on parsers and templates in the `contrib/` area of the PostgreSQL distribution.) Since dictionaries and configurations just parameterize and connect together some underlying parsers and templates, no special privilege is needed to create a new dictionary or configuration. Examples of creating custom dictionaries and configurations appear later in this chapter.

### 12.2. Tables and Indexes

The examples in the previous section illustrated full text matching using simple constant strings. This section shows how to search table data, optionally using indexes.

#### 12.2.1. Searching a Table

It is possible to do a full text search without an index. A simple query to print the `title` of each row that contains the word `friend` in its `body` field is:

```sql
SELECT title
FROM pgweb
WHERE to_tsvector('english', body) @@ to_tsquery('english', 'friend');
```

This will also find related words such as `friends` and `friendly`, since all these are reduced to the same normalized lexeme.

The query above specifies that the `english` configuration is to be used to parse and normalize the strings. Alternatively we could omit the configuration parameters:

```sql
SELECT title
FROM pgweb
WHERE to_tsvector(body) @@ to_tsquery('friend');
```

This query will use the configuration set by `default_text_search_config`.

A more complex example is to select the ten most recent documents that contain `create` and `table` in the `title` or `body`:

```sql
SELECT title
FROM pgweb
```
WHERE to_tsvector(title || ' ' || body) @@ to_tsquery('create & table')
ORDER BY last_mod_date DESC
LIMIT 10;

For clarity we omitted the coalesce function calls which would be needed to find rows that contain NULL in one of the two fields.

Although these queries will work without an index, most applications will find this approach too slow, except perhaps for occasional ad-hoc searches. Practical use of text searching usually requires creating an index.

### 12.2.2. Creating Indexes

We can create a GIN index (Section 12.9) to speed up text searches:

```sql
CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector('english', body));
```

Notice that the 2-argument version of `to_tsvector` is used. Only text search functions that specify a configuration name can be used in expression indexes (Section 11.7). This is because the index contents must be unaffected by default_text_search_config. If they were affected, the index contents might be inconsistent because different entries could contain `tsvector`s that were created with different text search configurations, and there would be no way to guess which was which. It would be impossible to dump and restore such an index correctly.

Because the two-argument version of `to_tsvector` was used in the index above, only a query reference that uses the 2-argument version of `to_tsvector` with the same configuration name will use that index. That is, WHERE `to_tsvector('english', body)` @@ 'a & b' can use the index, but WHERE `to_tsvector(body)` @@ 'a & b' cannot. This ensures that an index will be used only with the same configuration used to create the index entries.

It is possible to set up more complex expression indexes wherein the configuration name is specified by another column, e.g.:

```sql
CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector(config_name, body));
```

where `config_name` is a column in the `pgweb` table. This allows mixed configurations in the same index while recording which configuration was used for each index entry. This would be useful, for example, if the document collection contained documents in different languages. Again, queries that are meant to use the index must be phrased to match, e.g., WHERE `to_tsvector(config_name, body)` @@ 'a & b'.

Indexes can even concatenate columns:

```sql
CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector('english', title || ' ' || body));
```

Another approach is to create a separate `tsvector` column to hold the output of `to_tsvector`. This example is a concatenation of `title` and `body`, using `coalesce` to ensure that one field will still be indexed when the other is NULL:

```sql
ALTER TABLE pgweb ADD COLUMN textsearchable_index_col tsvector;
UPDATE pgweb SET textsearchable_index_col =
```
Then we create a GIN index to speed up the search:

```sql
CREATE INDEX textsearch_idx ON pgweb USING GIN
(textsearchable_index_col);
```

Now we are ready to perform a fast full text search:

```sql
SELECT title
FROM pgweb
WHERE textsearchable_index_col @@ to_tsquery('create & table')
ORDER BY last_mod_date DESC
LIMIT 10;
```

When using a separate column to store the tsvector representation, it is necessary to create a trigger to keep the tsvector column current anytime title or body changes. Section 12.4.3 explains how to do that.

One advantage of the separate-column approach over an expression index is that it is not necessary to explicitly specify the text search configuration in queries in order to make use of the index. As shown in the example above, the query can depend on default_text_search_config. Another advantage is that searches will be faster, since it will not be necessary to redo the to_tsvector calls to verify index matches. (This is more important when using a GiST index than a GIN index; see Section 12.9.) The expression-index approach is simpler to set up, however, and it requires less disk space since the tsvector representation is not stored explicitly.

## 12.3. Controlling Text Search

To implement full text searching there must be a function to create a tsvector from a document and a tsquery from a user query. Also, we need to return results in a useful order, so we need a function that compares documents with respect to their relevance to the query. It's also important to be able to display the results nicely. PostgreSQL provides support for all of these functions.

### 12.3.1. Parsing Documents

PostgreSQL provides the function `to_tsvector` for converting a document to the tsvector data type.

```sql
to_tsvector([ config regconfig, ] document text) returns tsvector
to_tsvector parses a textual document into tokens, reduces the tokens to lexemes, and returns a tsvector which lists the lexemes together with their positions in the document. The document is processed according to the specified or default text search configuration. Here is a simple example:

```
SELECT to_tsvector('english', 'a fat cat sat on a mat - it ate a fat rats');
```

```
to_tsvector('english', 'a fat cat sat on a mat - it ate a fat rats')
------------------------------
'ate':9 'cat':3 'fat':2,11 'mat':7 'rat':12 'sat':4
```

In the example above we see that the resulting tsvector does not contain the words a, on, or it, the word rats became rat, and the punctuation sign – was ignored.
The \texttt{to_tsvector} function internally calls a parser which breaks the document text into tokens and assigns a type to each token. For each token, a list of dictionaries (Section 12.6) is consulted, where the list can vary depending on the token type. The first dictionary that recognizes the token emits one or more normalized lexemes to represent the token. For example, \texttt{rats} became \texttt{rat} because one of the dictionaries recognized that the word \texttt{rats} is a plural form of \texttt{rat}. Some words are recognized as stop words (Section 12.6.1), which causes them to be ignored since they occur too frequently to be useful in searching. In our example these are \texttt{a}, \texttt{on}, and \texttt{it}. If no dictionary in the list recognizes the token then it is also ignored. In this example that happened to the punctuation sign - because there are in fact no dictionaries assigned for its token type (Space symbols), meaning space tokens will never be indexed. The choices of parser, dictionaries and which types of tokens to index are determined by the selected text search configuration (Section 12.7). It is possible to have many different configurations in the same database, and predefined configurations are available for various languages. In our example we used the default configuration \texttt{english} for the English language.

The function \texttt{setweight} can be used to label the entries of a \texttt{tsvector} with a given weight, where a weight is one of the letters \texttt{A}, \texttt{B}, \texttt{C}, or \texttt{D}. This is typically used to mark entries coming from different parts of a document, such as title versus body. Later, this information can be used for ranking of search results.

Because \texttt{to_tsvector(NULL)} will return NULL, it is recommended to use \texttt{coalesce} whenever a field might be null. Here is the recommended method for creating a \texttt{tsvector} from a structured document:

\begin{verbatim}
UPDATE tt SET ti =
    setweight(to_tsvector(coalesce(title,'')), 'A') ||
    setweight(to_tsvector(coalesce(keyword,'')), 'B') ||
    setweight(to_tsvector(coalesce(abstract,'')), 'C') ||
    setweight(to_tsvector(coalesce(body,'')), 'D');
\end{verbatim}

Here we have used \texttt{setweight} to label the source of each lexeme in the finished \texttt{tsvector}, and then merged the labeled \texttt{tsvector} values using the \texttt{tsvector} concatenation operator ||. (Section 12.4.1 gives details about these operations.)

### 12.3.2. Parsing Queries

PostgreSQL provides the functions \texttt{to_tquery}, \texttt{plainto_tquery}, \texttt{phraseto_tquery} and \texttt{websearch_to_tquery} for converting a query to the \texttt{tsquery} data type. \texttt{to_tquery} offers access to more features than either \texttt{plainto_tquery} or \texttt{phraseto_tquery}, but it is less forgiving about its input. \texttt{websearch_to_tquery} is a simplified version of \texttt{to_tquery} with an alternative syntax, similar to the one used by web search engines.

\begin{verbatim}
to_tquery([ config regconfig, ] querytext text) returns tsquery
\end{verbatim}

\texttt{to_tquery} creates a \texttt{tsquery} value from \texttt{querytext}, which must consist of single tokens separated by the \texttt{tsquery} operators \& (AND), \mid (OR), \! (NOT), and \texttt{<>} (FOLLOWED BY), possibly grouped using parentheses. In other words, the input to \texttt{to_tquery} must already follow the general rules for \texttt{tsquery} input, as described in Section 8.11.2. The difference is that while basic \texttt{tsquery} input takes the tokens at face value, \texttt{to_tquery} normalizes each token into a lexeme using the specified or default configuration, and discards any tokens that are stop words according to the configuration. For example:

\begin{verbatim}
SELECT to_tquery('english', 'The & Fat & Rats');
\end{verbatim}

---

392
As in basic \texttt{tsquery} input, weight(s) can be attached to each lexeme to restrict it to match only \texttt{tsvector} lexemes of those weight(s). For example:

\begin{verbatim}
SELECT to_tsquery('english', 'Fat | Rats:AB');
  to_tsquery
  ------------------
    'fat' | 'rat':AB
\end{verbatim}

Also, * can be attached to a lexeme to specify prefix matching:

\begin{verbatim}
SELECT to_tsquery('supern:*A & star:A*B');
  to_tsquery
  --------------------------
    'supern':*A & 'star':*AB
\end{verbatim}

Such a lexeme will match any word in a \texttt{tsvector} that begins with the given string.

to\_tsquery can also accept single-quoted phrases. This is primarily useful when the configuration includes a thesaurus dictionary that may trigger on such phrases. In the example below, a thesaurus contains the rule supernovae stars : sn:

\begin{verbatim}
SELECT to_tsquery('''supernovae stars'' & !crab');
  to_tsquery
  -----------------
    'sn' & '!crab'
\end{verbatim}

Without quotes, \texttt{to\_tsquery} will generate a syntax error for tokens that are not separated by an AND, OR, or FOLLOWED BY operator.

\begin{verbatim}
plainto\_tsquery([ config regconfig, ] querytext text)
  returns tsquery

plainto\_tsquery transforms the unformatted text querytext to a tsquery value. The text is parsed and normalized much as for to\_tsvector, then the & (AND) tsquery operator is inserted between surviving words.

Example:

\begin{verbatim}
SELECT plainto\_tsquery('english', 'The Fat Rats');
  plainto\_tsquery
  -----------------
    'fat' & 'rat'
\end{verbatim}

Note that plainto\_tsquery will not recognize tsquery operators, weight labels, or prefix-match labels in its input:

\begin{verbatim}
SELECT plainto\_tsquery('english', 'The Fat & Rats:C');
  plainto\_tsquery
  -----------------
    'fat' & 'rat' & 'c'
\end{verbatim}

Here, all the input punctuation was discarded as being space symbols.
phraseto_tsquery([ config regconfig, ] querytext text) returns tsquery

phraseto_tsquery behaves much like plainto_tsquery, except that it inserts the <-> (FOLLOWED BY) operator between surviving words instead of the & (AND) operator. Also, stop words are not simply discarded, but are accounted for by inserting <N> operators rather than <-> operators. This function is useful when searching for exact lexeme sequences, since the FOLLOWED BY operators check lexeme order not just the presence of all the lexemes.

Example:

SELECT phraseto_tsquery('english', 'The Fat Rats');
phraseto_tsquery
------------------
'fat' <-> 'rat'

Like plainto_tsquery, the phraseto_tsquery function will not recognize tsquery operators, weight labels, or prefix-match labels in its input:

SELECT phraseto_tsquery('english', 'The Fat & Rats:C');
phraseto_tsquery
------------------
'fat' <-> 'rat' <-> 'c'

websearch_to_tsquery([ config regconfig, ] querytext text) returns tsquery

websearch_to_tsquery creates a tsquery value from querytext using an alternative syntax in which simple unformatted text is a valid query. Unlike plainto_tsquery and phraseto_tsquery, it also recognizes certain operators. Moreover, this function should never raise syntax errors, which makes it possible to use raw user-supplied input for search. The following syntax is supported:

• unquoted text: text not inside quote marks will be converted to terms separated by & operators, as if processed by plainto_tsquery.
• "quoted text": text inside quote marks will be converted to terms separated by <-> operators, as if processed by phraseto_tsquery.
• OR: logical or will be converted to the | operator.
• -: the logical not operator, converted to the the ! operator.

Examples:

SELECT websearch_to_tsquery('english', 'The fat rats');
websearch_to_tsquery
------------------
'fat' & 'rat'
(1 row)

SELECT websearch_to_tsquery('english', '"supernovae stars" -crab');
websearch_to_tsquery
------------------
'supernova' <-> 'star' & !'crab'
(1 row)

SELECT websearch_to_tsquery('english', '"sad cat" or "fat rat"');
websearch_to_tsquery
------------------
12.3.3. Ranking Search Results

Ranking attempts to measure how relevant documents are to a particular query, so that when there are many matches the most relevant ones can be shown first. PostgreSQL provides two predefined ranking functions, which take into account lexical, proximity, and structural information; that is, they consider how often the query terms appear in the document, how close together the terms are in the document, and how important is the part of the document where they occur. However, the concept of relevancy is vague and very application-specific. Different applications might require additional information for ranking, e.g., document modification time. The built-in ranking functions are only examples. You can write your own ranking functions and/or combine their results with additional factors to fit your specific needs.

The two ranking functions currently available are:

- `ts_rank(weights float4[], vector tsvector, query tsquery [, normalization integer ]) returns float4`
  Ranks vectors based on the frequency of their matching lexemes.

- `ts_rank_cd(weights float4[], vector tsvector, query tsquery [, normalization integer ]) returns float4`
  This function computes the cover density ranking for the given document vector and query, as described in Clarke, Cormack, and Tudhope's "Relevance Ranking for One to Three Term Queries" in the journal "Information Processing and Management", 1999. Cover density is similar to `ts_rank` ranking except that the proximity of matching lexemes to each other is taken into consideration.

  This function requires lexeme positional information to perform its calculation. Therefore, it ignores any "stripped" lexemes in the `tsvector`. If there are no unstripped lexemes in the input, the result will be zero. (See Section 12.4.1 for more information about the `strip` function and positional information in `tsvector`.)

For both these functions, the optional `weights` argument offers the ability to weigh word instances more or less heavily depending on how they are labeled. The weight arrays specify how heavily to weigh each category of word, in the order:

- `D-weight`
- `C-weight`
- `B-weight`
- `A-weight`

If no `weights` are provided, then these defaults are used:
Typically weights are used to mark words from special areas of the document, like the title or an initial abstract, so they can be treated with more or less importance than words in the document body.

Since a longer document has a greater chance of containing a query term it is reasonable to take into account document size, e.g., a hundred-word document with five instances of a search word is probably more relevant than a thousand-word document with five instances. Both ranking functions take an integer normalization option that specifies whether and how a document’s length should impact its rank. The integer option controls several behaviors, so it is a bit mask: you can specify one or more behaviors using | (for example, 2 | 4).

- 0 (the default) ignores the document length
- 1 divides the rank by 1 + the logarithm of the document length
- 2 divides the rank by the document length
- 4 divides the rank by the mean harmonic distance between extents (this is implemented only by ts_rank_cd)
- 8 divides the rank by the number of unique words in document
- 16 divides the rank by 1 + the logarithm of the number of unique words in document
- 32 divides the rank by itself + 1

If more than one flag bit is specified, the transformations are applied in the order listed.

It is important to note that the ranking functions do not use any global information, so it is impossible to produce a fair normalization to 1% or 100% as sometimes desired. Normalization option 32 (rank/(rank+1)) can be applied to scale all ranks into the range zero to one, but of course this is just a cosmetic change; it will not affect the ordering of the search results.

Here is an example that selects only the ten highest-ranked matches:

```
SELECT title, ts_rank_cd(textsearch, query) AS rank
FROM apod, to_tsquery('neutrino| (dark & matter)') query
WHERE query @@ textsearch
ORDER BY rank DESC
LIMIT 10;
```

<table>
<thead>
<tr>
<th>title</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>Neutrinos in the Sun</td>
<td>3.1</td>
</tr>
<tr>
<td>The Sudbury Neutrino Detector</td>
<td>2.4</td>
</tr>
<tr>
<td>A MACHO View of Galactic Dark Matter</td>
<td>2.01317</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.91171</td>
</tr>
<tr>
<td>The Virgo Cluster: Hot Plasma and Dark Matter</td>
<td>1.90953</td>
</tr>
<tr>
<td>Rafting for Solar Neutrinos</td>
<td>1.9</td>
</tr>
<tr>
<td>NGC 4650A: Strange Galaxy and Dark Matter</td>
<td>1.85774</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.6123</td>
</tr>
<tr>
<td>Ice Fishing for Cosmic Neutrinos</td>
<td>1.6</td>
</tr>
<tr>
<td>Weak Lensing Distorts the Universe</td>
<td>0.818218</td>
</tr>
</tbody>
</table>

This is the same example using normalized ranking:

```
SELECT title, ts_rank_cd(textsearch, query, 32 /* rank/(rank+1) */ ) AS rank
FROM apod, to_tsquery('neutrino| (dark & matter)') query
WHERE query @@ textsearch
ORDER BY rank DESC
LIMIT 10;
```

<table>
<thead>
<tr>
<th>title</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>Neutrinos in the Sun</td>
<td>3.1</td>
</tr>
<tr>
<td>The Sudbury Neutrino Detector</td>
<td>2.4</td>
</tr>
<tr>
<td>A MACHO View of Galactic Dark Matter</td>
<td>2.01317</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.91171</td>
</tr>
<tr>
<td>The Virgo Cluster: Hot Plasma and Dark Matter</td>
<td>1.90953</td>
</tr>
<tr>
<td>Rafting for Solar Neutrinos</td>
<td>1.9</td>
</tr>
<tr>
<td>NGC 4650A: Strange Galaxy and Dark Matter</td>
<td>1.85774</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.6123</td>
</tr>
<tr>
<td>Ice Fishing for Cosmic Neutrinos</td>
<td>1.6</td>
</tr>
<tr>
<td>Weak Lensing Distorts the Universe</td>
<td>0.818218</td>
</tr>
</tbody>
</table>
Ranking can be expensive since it requires consulting the tsvector of each matching document, which can be I/O bound and therefore slow. Unfortunately, it is almost impossible to avoid since practical queries often result in large numbers of matches.

### 12.3.4. Highlighting Results

To present search results it is ideal to show a part of each document and how it is related to the query. Usually, search engines show fragments of the document with marked search terms. PostgreSQL provides a function `ts_headline` that implements this functionality.

```sql
ts_headline([ config regconfig, ] document text, query tsquery [, options text ]) returns text
```

`ts_headline` accepts a document along with a query, and returns an excerpt from the document in which terms from the query are highlighted. The configuration to be used to parse the document can be specified by `config`; if `config` is omitted, the `default_text_search_config` configuration is used.

If an `options` string is specified it must consist of a comma-separated list of one or more `option=value` pairs. The available options are:

- **StartSel, StopSel**: the strings with which to delimit query words appearing in the document, to distinguish them from other excerpted words. You must double-quote these strings if they contain spaces or commas.
- **MaxWords, MinWords**: these numbers determine the longest and shortest headlines to output.
- **ShortWord**: words of this length or less will be dropped at the start and end of a headline. The default value of three eliminates common English articles.
- **HighlightAll**: Boolean flag; if `true` the whole document will be used as the headline, ignoring the preceding three parameters.
- **MaxFragments**: maximum number of text excerpts or fragments to display. The default value of zero selects a non-fragment-oriented headline generation method. A value greater than zero selects fragment-based headline generation. This method finds text fragments with as many query words as possible and stretches those fragments around the query words. As a result query words are close to the middle of each fragment and have words on each side. Each fragment will be of at most `MaxWords` and words of length `ShortWord` or less are dropped at the start and end of each fragment. If not all query words are found in the document, then a single fragment of the first `MinWords` in the document will be displayed.
- **FragmentDelimiter**: When more than one fragment is displayed, the fragments will be separated by this string.

These option names are recognized case-insensitively. Any unspecified options receive these defaults:

```
StartSel=<b>, StopSel=</b>,
MaxWords=35, MinWords=15, ShortWord=3, HighlightAll=FALSE,
```
MaxFragments=0, FragmentDelimiter="..."

For example:

```sql
SELECT ts_headline('english',
   'The most common type of search
is to find all documents containing given query terms
and return them in order of their similarity to the
query.',
   to_tsquery('query & similarity'));
```

For example:

```sql
SELECT ts_headline('english',
   'The most common type of search
is to find all documents containing given query terms
and return them in order of their similarity to the
query.',
   to_tsquery('query & similarity'),
   'StartSel = <, StopSel = >');
```

12.4. Additional Features

This section describes additional functions and operators that are useful in connection with text search.

12.4.1. Manipulating Documents

Section 12.3.1 showed how raw textual documents can be converted into `tsvector` values. PostgreSQL also provides functions and operators that can be used to manipulate documents that are already in `tsvector` form.

```sql
tsvector || tsvector
```

The `tsvector` concatenation operator returns a vector which combines the lexemes and positional information of the two vectors given as arguments. Positions and weight labels are retained during the concatenation. Positions appearing in the right-hand vector are offset by the largest position mentioned in the left-hand vector, so that the result is nearly equivalent to the result of performing `to_tsvector` on the concatenation of the two original document strings. (The equivalence is not exact, because any stop-words removed from the end of the left-hand argument will not affect the result, whereas they would have affected the positions of the lexemes in the right-hand argument if textual concatenation were used.)

One advantage of using concatenation in the vector form, rather than concatenating text before applying `to_tsvector`, is that you can use different configurations to parse different sections of the document. Also, because the `setweight` function marks all lexemes of the given vector the same way, it is necessary to parse the text and do `setweight` before concatenating if you want to label different parts of the document with different weights.
setweight(vector tsvector, weight "char") returns tsvector

setweight returns a copy of the input vector in which every position has been labeled with the given weight, either A, B, C, or D. (D is the default for new vectors and as such is not displayed on output.) These labels are retained when vectors are concatenated, allowing words from different parts of a document to be weighted differently by ranking functions.

Note that weight labels apply to positions, not lexemes. If the input vector has been stripped of positions then setweight does nothing.

length(vector tsvector) returns integer

Returns the number of lexemes stored in the vector.

strip(vector tsvector) returns tsvector

Returns a vector that lists the same lexemes as the given vector, but lacks any position or weight information. The result is usually much smaller than an unstripped vector, but it is also less useful. Relevance ranking does not work as well on stripped vectors as unstripped ones. Also, the <-> (FOLLOWED BY) tsquery operator will never match stripped input, since it cannot determine the distance between lexeme occurrences.

A full list of tsvector-related functions is available in Table 9.41.

### 12.4.2. Manipulating Queries

Section 12.3.2 showed how raw textual queries can be converted into tsquery values. PostgreSQL also provides functions and operators that can be used to manipulate queries that are already in tsquery form.

**tsquery & tsquery**

Returns the AND-combination of the two given queries.

**tsquery | tsquery**

Returns the OR-combination of the two given queries.

**!! tsquery**

Returns the negation (NOT) of the given query.

**tsquery <-> tsquery**

Returns a query that searches for a match for the first given query immediately followed by a match to the second given query, using the <-> (FOLLOWED BY) tsquery operator. For example:

```sql
SELECT to_tsquery('fat') <-> to_tsquery('cat | rat');
```

**tsquery_phrase(query1 tsquery, query2 tsquery [, distance integer]) returns tsquery**

Returns a query that searches for a match to the first given query followed by a match to the second given query at a distance of at distance lexemes, using the <N> tsquery operator. For example:

```sql
SELECT tsquery_phrase(to_tsquery('fat'), to_tsquery('cat'), 10); tsquery_phrase
```
'fat' <10> 'cat'

\texttt{numnode(query \texttt{tsquery}) \textbf{returns integer}}

Returns the number of nodes (lexemes plus operators) in a \texttt{tsquery}. This function is useful to determine if the \texttt{query} is meaningful (returns > 0), or contains only stop words (returns 0).

\textbf{Examples:}

\begin{verbatim}
SELECT numnode(plainto_tsquery('the any'));
\end{verbatim}

\texttt{NOTICE: query contains only stopword(s) or doesn't contain lexeme(s), ignored}

\texttt{numnode} ---------

\texttt{0}

\begin{verbatim}
SELECT numnode('foo \& bar':tsquery);
\end{verbatim}

\texttt{numnode} ---------

\texttt{3}

\texttt{querytree(query \texttt{tsquery}) \textbf{returns text}}

Returns the portion of a \texttt{tsquery} that can be used for searching an index. This function is useful for detecting unindexable queries, for example those containing only stop words or only negated terms. For example:

\begin{verbatim}
SELECT querytree(to_tsquery('!defined'));
\end{verbatim}

\texttt{querytree} -----------

\textbf{12.4.2.1. Query Rewriting}

The \texttt{ts_rewrite} family of functions search a given \texttt{tsquery} for occurrences of a target subquery, and replace each occurrence with a substitute subquery. In essence this operation is a \texttt{tsquery}-specific version of substring replacement. A target and substitute combination can be thought of as a \texttt{query rewrite rule}. A collection of such rewrite rules can be a powerful search aid. For example, you can expand the search using synonyms (e.g., \texttt{new york}, \texttt{big apple}, \texttt{nyc}, \texttt{gotham}) or narrow the search to direct the user to some hot topic. There is some overlap in functionality between this feature and thesaurus dictionaries (Section 12.6.4). However, you can modify a set of rewrite rules on-the-fly without reindexing, whereas updating a thesaurus requires reindexing to be effective.

\texttt{ts_rewrite (query \texttt{tsquery}, target \texttt{tsquery}, substitute \texttt{tsquery}) \textbf{returns tsquery}}

This form of \texttt{ts_rewrite} simply applies a single rewrite rule: \texttt{target} is replaced by \texttt{substitute} wherever it appears in \texttt{query}. For example:

\begin{verbatim}
SELECT ts_rewrite('a \& b':tsquery, 'a':tsquery, 'c':tsquery);
\end{verbatim}

\texttt{ts_rewrite} ---------

\texttt{'b' \& 'c'}

\texttt{ts_rewrite (query \texttt{tsquery}, select \texttt{text}) \textbf{returns tsquery}}

This form of \texttt{ts_rewrite} accepts a starting \texttt{query} and a SQL \texttt{select} command, which is given as a text string. The \texttt{select} must yield two columns of \texttt{tsquery} type. For each row of
the select result, occurrences of the first column value (the target) are replaced by the second column value (the substitute) within the current query value. For example:

```sql
CREATE TABLE aliases (t tsquery PRIMARY KEY, s tsquery);
INSERT INTO aliases VALUES('a', 'c');

SELECT ts_rewrite('a & b '::tsquery, 'SELECT t, s FROM aliases');
```

Note that when multiple rewrite rules are applied in this way, the order of application can be important; so in practice you will want the source query to ORDER BY some ordering key.

Let's consider a real-life astronomical example. We'll expand query supernovae using table-driven rewriting rules:

```sql
CREATE TABLE aliases (t tsquery primary key, s tsquery);
INSERT INTO aliases VALUES(to_tsquery('supernovae'),
                           to_tsquery('supernovae|sn'));

SELECT ts_rewrite(to_tsquery('supernovae & crab'), 'SELECT * FROM aliases');
```

We can change the rewriting rules just by updating the table:

```sql
UPDATE aliases
SET s = to_tsquery('supernovae|sn & !nebulae')
WHERE t = to_tsquery('supernovae');

SELECT ts_rewrite(to_tsquery('supernovae & crab'), 'SELECT * FROM aliases');
```

Rewriting can be slow when there are many rewriting rules, since it checks every rule for a possible match. To filter out obvious non-candidate rules we can use the containment operators for the tsquery type. In the example below, we select only those rules which might match the original query:

```sql
SELECT ts_rewrite('a & b '::tsquery, 
                 'SELECT t, s FROM aliases WHERE ''a & b'' '::tsquery @> t');
```

12.4.3. Triggers for Automatic Updates

When using a separate column to store the tsvector representation of your documents, it is necessary to create a trigger to update the tsvector column when the document content columns change. Two built-in trigger functions are available for this, or you can write your own.
These trigger functions automatically compute a \texttt{tsvector} column from one or more textual columns, under the control of parameters specified in the \texttt{CREATE TRIGGER} command. An example of their use is:

\begin{verbatim}
CREATE TABLE messages (
  title       text,
  body        text,
  tsv         tsvector
);

CREATE TRIGGER tsvectorupdate BEFORE INSERT OR UPDATE
ON messages FOR EACH ROW EXECUTE FUNCTION
  tsvector_update_trigger(tsv, 'pg_catalog.english', title, body);

INSERT INTO messages VALUES('title here', 'the body text is here');

SELECT * FROM messages;
  title    |         body          |            tsv
------------+-----------------------+----------------------------
  title here | the body text is here | 'bodi':4 'text':5 'titl':1

SELECT title, body FROM messages WHERE tsv @@ to_tsquery('title & body');
  title    |         body
------------+-----------------------
  title here | the body text is here

Having created this trigger, any change in \texttt{title} or \texttt{body} will automatically be reflected into \texttt{tsv}, without the application having to worry about it.

The first trigger argument must be the name of the \texttt{tsvector} column to be updated. The second argument specifies the text search configuration to be used to perform the conversion. For \texttt{tsvector_update_trigger}, the configuration name is simply given as the second trigger argument. It must be schema-qualified as shown above, so that the trigger behavior will not change with changes in \texttt{search_path}. For \texttt{tsvector_update_trigger_column}, the second trigger argument is the name of another table column, which must be of type \texttt{regconfig}. This allows a per-row selection of configuration to be made. The remaining argument(s) are the names of textual columns (of type \texttt{text}, \texttt{varchar}, or \texttt{char}). NULL values will be skipped (but the other columns will still be indexed).

A limitation of these built-in triggers is that they treat all the input columns alike. To process columns differently — for example, to weight title differently from body — it is necessary to write a custom trigger. Here is an example using \texttt{PL/pgSQL} as the trigger language:

\begin{verbatim}
CREATE FUNCTION messages_trigger() RETURNS trigger AS $$
begin
  new.tsv :=
    setweight(to_tsvector('pg_catalog.english',
      coalesce(new.title,'')), 'A') ||
    setweight(to_tsvector('pg_catalog.english',
      coalesce(new.body,'')), 'D');
end;
$$
\end{verbatim}
12.4.4. Gathering Document Statistics

The function `ts_stat` is useful for checking your configuration and for finding stop-word candidates.

\[
\text{ts_stat(sqlquery text, [ weights text, ]}
\text{ OUT word text, OUT ndoc integer,}
\text{ OUT nentry integer) returns setof record}
\]

*sqlquery* is a text value containing an SQL query which must return a single `tsvector` column. `ts_stat` executes the query and returns statistics about each distinct lexeme (word) contained in the `tsvector` data. The columns returned are

- *word* text — the value of a lexeme
- *ndoc* integer — number of documents (`tsvector`s) the word occurred in
- *nentry* integer — total number of occurrences of the word

If *weights* is supplied, only occurrences having one of those weights are counted.

For example, to find the ten most frequent words in a document collection:

\[
\text{SELECT * FROM ts_stat('SELECT vector FROM apod')}
\text{ORDER BY nentry DESC, ndoc DESC, word}
\text{LIMIT 10;}
\]

The same, but counting only word occurrences with weight A or B:

\[
\text{SELECT * FROM ts_stat('SELECT vector FROM apod', 'ab')}
\text{ORDER BY nentry DESC, ndoc DESC, word}
\text{LIMIT 10;}
\]

12.5. Parsers

Text search parsers are responsible for splitting raw document text into tokens and identifying each token’s type, where the set of possible types is defined by the parser itself. Note that a parser does not modify the text at all — it simply identifies plausible word boundaries. Because of this limited scope, there is less need for application-specific custom parsers than there is for custom dictionaries. At present PostgreSQL provides just one built-in parser, which has been found to be useful for a wide range of applications.

The built-in parser is named `pg_catalog.default`. It recognizes 23 token types, shown in Table 12.1.
Table 12.1. Default Parser's Token Types

<table>
<thead>
<tr>
<th>Alias</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII letters</td>
<td>elephant</td>
</tr>
<tr>
<td>word</td>
<td>Word, all letters</td>
<td>mañana</td>
</tr>
<tr>
<td>numword</td>
<td>Word, letters and digits</td>
<td>beta1</td>
</tr>
<tr>
<td>asciihword</td>
<td>Hyphenated word, all ASCII</td>
<td>up-to-date</td>
</tr>
<tr>
<td>hword</td>
<td>Hyphenated word, all letters</td>
<td>lógico-matemática</td>
</tr>
<tr>
<td>numhword</td>
<td>Hyphenated word, letters and digits</td>
<td>postgresql-beta1</td>
</tr>
<tr>
<td>hword_asciipart</td>
<td>Hyphenated word part, all ASCII</td>
<td>postgresql in the context</td>
</tr>
<tr>
<td>hword_part</td>
<td>Hyphenated word part, all letters</td>
<td>lógico or matemática in the context</td>
</tr>
<tr>
<td>hword_numpart</td>
<td>Hyphenated word part, letters and digits</td>
<td>betal in the context</td>
</tr>
<tr>
<td>email</td>
<td>Email address</td>
<td><a href="mailto:foo@example.com">foo@example.com</a></td>
</tr>
<tr>
<td>protocol</td>
<td>Protocol head</td>
<td>http://</td>
</tr>
<tr>
<td>url</td>
<td>URL</td>
<td>example.com/stuff/index.html</td>
</tr>
<tr>
<td>host</td>
<td>Host</td>
<td>example.com</td>
</tr>
<tr>
<td>url_path</td>
<td>URL path</td>
<td>/stuff/index.html, in the context of a URL</td>
</tr>
<tr>
<td>file</td>
<td>File or path name</td>
<td>/usr/local/foo.txt, if not within a URL</td>
</tr>
<tr>
<td>sfloat</td>
<td>Scientific notation</td>
<td>-1.234e56</td>
</tr>
<tr>
<td>float</td>
<td>Decimal notation</td>
<td>-1.234</td>
</tr>
<tr>
<td>int</td>
<td>Signed integer</td>
<td>-1234</td>
</tr>
<tr>
<td>uint</td>
<td>Unsigned integer</td>
<td>1234</td>
</tr>
<tr>
<td>version</td>
<td>Version number</td>
<td>8.3.0</td>
</tr>
<tr>
<td>tag</td>
<td>XML tag</td>
<td>&lt;a href=&quot;dictionaries.html&quot;&gt;</td>
</tr>
<tr>
<td>entity</td>
<td>XML entity</td>
<td>&amp;</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td>(any whitespace or punctuation not otherwise recognized)</td>
</tr>
</tbody>
</table>

**Note**

The parser's notion of a “letter” is determined by the database’s locale setting, specifically lc_ctype. Words containing only the basic ASCII letters are reported as a separate token type, since it is sometimes useful to distinguish them. In most European languages, token types word and asciiword should be treated alike.

email does not support all valid email characters as defined by RFC 5322. Specifically, the only non-alphanumeric characters supported for email user names are period, dash, and underscore.
It is possible for the parser to produce overlapping tokens from the same piece of text. As an example, a hyphenated word will be reported both as the entire word and as each component:

```
SELECT alias, description, token FROM ts_debug('foo-bar-beta1');
alias |               description                |
-----------------+------------------------------------------
numhword        | Hyphenated word, letters and digits      | foo-
bar-beta1       |                                          |
```

This behavior is desirable since it allows searches to work for both the whole compound word and for components. Here is another instructive example:

```
SELECT alias, description, token FROM ts_debug('http://example.com/stuff/index.html');
alias   |  description  |            token
----------+---------------+-------------------------------
protocol | Protocol head | http://
url      | URL           | example.com/stuff/index.html
host     | Host          | example.com
url_path | URL path      | /stuff/index.html
```

### 12.6. Dictionaries

Dictionaries are used to eliminate words that should not be considered in a search (*stop words*), and to normalize words so that different derived forms of the same word will match. A successfully normalized word is called a *lexeme*. Aside from improving search quality, normalization and removal of stop words reduce the size of the *tsvector* representation of a document, thereby improving performance. Normalization does not always have linguistic meaning and usually depends on application semantics.

Some examples of normalization:

- **Linguistic** - Ispell dictionaries try to reduce input words to a normalized form; stemmer dictionaries remove word endings
- **URL locations** can be canonicalized to make equivalent URLs match:
  - `http://www.pgsq1.ru/db/mw/`
- **Color names** can be replaced by their hexadecimal values, e.g., `red`, `green`, `blue`, `magenta` -> `FF0000`, `00FF00`, `0000FF`, `FF00FF`
- If indexing numbers, we can remove some fractional digits to reduce the range of possible numbers, so for example `3.14159265359, 3.1415926, 3.14` will be the same after normalization if only two digits are kept after the decimal point.

A dictionary is a program that accepts a token as input and returns:

- an array of lexemes if the input token is known to the dictionary (notice that one token can produce more than one lexeme)
- a single lexeme with the *TSL_FILTER* flag set, to replace the original token with a new token to be passed to subsequent dictionaries (a dictionary that does this is called a *filtering dictionary*)
Full Text Search

- an empty array if the dictionary knows the token, but it is a stop word
- NULL if the dictionary does not recognize the input token

PostgreSQL provides predefined dictionaries for many languages. There are also several predefined
templates that can be used to create new dictionaries with custom parameters. Each predefined dictio-
nary template is described below. If no existing template is suitable, it is possible to create new ones;
see the contrib/ area of the PostgreSQL distribution for examples.

A text search configuration binds a parser together with a set of dictionaries to process the parser’s
output tokens. For each token type that the parser can return, a separate list of dictionaries is specified
by the configuration. When a token of that type is found by the parser, each dictionary in the list is
consulted in turn, until some dictionary recognizes it as a known word. If it is identified as a stop
word, or if no dictionary recognizes the token, it will be discarded and not indexed or searched for.
Normally, the first dictionary that returns a non-NULL output determines the result, and any remaining
dictionaries are not consulted; but a filtering dictionary can replace the given word with a modified
word, which is then passed to subsequent dictionaries.

The general rule for configuring a list of dictionaries is to place first the most narrow, most specific
dictionary, then the more general dictionaries, finishing with a very general dictionary, like a Snowball
stemmer or simple, which recognizes everything. For example, for an astronomy-specific search
(astro_en configuration) one could bind token type asciiword (ASCII word) to a synonym
dictionary of astronomical terms, a general English dictionary and a Snowball English stemmer:

```sql
ALTER TEXT SEARCH CONFIGURATION astro_en
    ADD MAPPING FOR asciiword WITH astrosyn, english_ispell, english_stem;
```

A filtering dictionary can be placed anywhere in the list, except at the end where it'd be useless.
Filtering dictionaries are useful to partially normalize words to simplify the task of later dictionaries.
For example, a filtering dictionary could be used to remove accents from accented letters, as is done
by the unaccent module.

### 12.6.1. Stop Words

Stop words are words that are very common, appear in almost every document, and have no discrim-
ination value. Therefore, they can be ignored in the context of full text searching. For example, every
English text contains words like *a* and *the*, so it is useless to store them in an index. However, stop
words do affect the positions in `tsvector`, which in turn affect ranking:

```sql
SELECT to_tsvector('english','in the list of stop words');
          to_tsvector
-----------------------------------------------
   'list':3 'stop':5 'word':6
```

The missing positions 1,2,4 are because of stop words. Ranks calculated for documents with and
without stop words are quite different:

```sql
SELECT ts_rank_cd (to_tsvector('english','in the list of stop words'), to_tsquery('list & stop'));
    ts_rank_cd
----------
      0.05
SELECT ts_rank_cd (to_tsvector('english','list stop words'), to_tsquery('list & stop'));
    ts_rank_cd
```
It is up to the specific dictionary how it treats stop words. For example, `ispell` dictionaries first normalize words and then look at the list of stop words, while `Snowball` stemmers first check the list of stop words. The reason for the different behavior is an attempt to decrease noise.

### 12.6.2. Simple Dictionary

The `simple` dictionary template operates by converting the input token to lower case and checking it against a file of stop words. If it is found in the file then an empty array is returned, causing the token to be discarded. If not, the lower-cased form of the word is returned as the normalized lexeme. Alternatively, the dictionary can be configured to report non-stop-words as unrecognized, allowing them to be passed on to the next dictionary in the list.

Here is an example of a dictionary definition using the `simple` template:

```sql
CREATE TEXT SEARCH DICTIONARY public.simple_dict (  
    TEMPLATE = pg_catalog.simple,  
    STOPWORDS = english  
);
```

Here, `english` is the base name of a file of stop words. The file's full name will be `$SHAREDIR/tsearch_data/english.stop`, where `$SHAREDIR` means the PostgreSQL installation's shared-data directory, often `/usr/local/share/postgresql` (use `pg_config --sharedir` to determine it if you're not sure). The file format is simply a list of words, one per line. Blank lines and trailing spaces are ignored, and upper case is folded to lower case, but no other processing is done on the file contents.

Now we can test our dictionary:

```sql
SELECT ts_lexize('public.simple_dict','YeS');
```

```
{yes}
```

```sql
SELECT ts_lexize('public.simple_dict','The');
```

```
{}
```

We can also choose to return `NULL`, instead of the lower-cased word, if it is not found in the stop words file. This behavior is selected by setting the dictionary's `Accept` parameter to `false`. Continuing the example:

```sql
ALTER TEXT SEARCH DICTIONARY public.simple_dict ( Accept = false );
```

```sql
SELECT ts_lexize('public.simple_dict','YeS');
```

```
--lexize
{yes}
```

```sql
SELECT ts_lexize('public.simple_dict','The');
```

```
{}
```
With the default setting of \texttt{Accept = true}, it is only useful to place a simple dictionary at the end of a list of dictionaries, since it will never pass on any token to a following dictionary. Conversely, \texttt{Accept = false} is only useful when there is at least one following dictionary.

\begin{center}
\underline{Caution}
\end{center}

Most types of dictionaries rely on configuration files, such as files of stop words. These files must be stored in UTF-8 encoding. They will be translated to the actual database encoding, if that is different, when they are read into the server.

\begin{center}
\underline{Caution}
\end{center}

Normally, a database session will read a dictionary configuration file only once, when it is first used within the session. If you modify a configuration file and want to force existing sessions to pick up the new contents, issue an \texttt{ALTER TEXT SEARCH DICTIONARY} command on the dictionary. This can be a “dummy” update that doesn’t actually change any parameter values.

### 12.6.3. Synonym Dictionary

This dictionary template is used to create dictionaries that replace a word with a synonym. Phrases are not supported (use the thesaurus template (Section 12.6.4) for that). A synonym dictionary can be used to overcome linguistic problems, for example, to prevent an English stemmer dictionary from reducing the word “Paris” to “pari”. It is enough to have a \texttt{Paris paris} line in the synonym dictionary and put it before the \texttt{english_stem} dictionary. For example:

```
SELECT * FROM ts_debug('english', 'Paris');
alias   |   description   | token |  dictionaries  |  dictionary
-----------+-----------------+-------+----------------
           | lexemes
+--------------+---------
 asciiword | Word, all ASCII | Paris | {english_stem} | {pari}
 english_stem | {pari}

CREATE TEXT SEARCH DICTIONARY my_synonym ( 
    TEMPLATE = synonym,
    SYNONYMS = my_synonyms
 );

ALTER TEXT SEARCH CONFIGURATION english
    ALTER MAPPING FOR asciiword
        WITH my_synonym, english_stem;

SELECT * FROM ts_debug('english', 'Paris');
alias   |   description   | token |       dictionaries        |
dictionary | lexemes |
-----------+-----------------+-------+---------------------------
           | lexemes
+--------------+---------
 asciiword | Word, all ASCII | Paris | {my_synonym,english_stem} |
 my_synonym   | {paris}

The only parameter required by the \texttt{synonym} template is \texttt{SYNONYMS}, which is the base name of its configuration file — \texttt{my_synonyms} in the above example. The file’s full name will be $SHAREDIR/
The file format is just one line per word to be substituted, with the word followed by its synonym, separated by white space. Blank lines and trailing spaces are ignored.

The synonym template also has an optional parameter CaseSensitive, which defaults to false. When CaseSensitive is false, words in the synonym file are folded to lower case, as are input tokens. When it is true, words and tokens are not folded to lower case, but are compared as-is.

An asterisk (*) can be placed at the end of a synonym in the configuration file. This indicates that the synonym is a prefix. The asterisk is ignored when the entry is used in to_tsvector(), but when it is used in to_tsquery(), the result will be a query item with the prefix match marker (see Section 12.3.2). For example, suppose we have these entries in $SHAREDIR/tsearch_data/synonym_sample.syn:

```
postgres       psql
postgresql     psql
postgresql     psql
postgresql     psql
indices        index*
```

Then we will get these results:

```
mydb=# CREATE TEXT SEARCH DICTIONARY syn (template=synonym,
synonyms='synonym_sample');
mydb=# SELECT ts_lexize('syn','indices');
t_lexize
  {index}
 (1 row)
mydb=# CREATE TEXT SEARCH CONFIGURATION tst (copy=simple);
mydb=# ALTER TEXT SEARCH CONFIGURATION tst ALTER MAPPING FOR
asciiword WITH syn;
mydb=# SELECT to_tsvector('tst','indices');
t_tsvector
  'index':1
 (1 row)
mydb=# SELECT to_tsquery('tst','indices');
t_tsquery
  'index':*
 (1 row)
mydb=# SELECT 'indexes are very useful':tsvector;
tvector
  'are' 'indexes' 'useful' 'very'
 (1 row)
mydb=# SELECT 'indexes are very useful':tsvector @@
to_tsquery('tst','indices');
?column?
 t
 (1 row)
```
12.6.4. Thesaurus Dictionary

A thesaurus dictionary (sometimes abbreviated as TZ) is a collection of words that includes information about the relationships of words and phrases, i.e., broader terms (BT), narrower terms (NT), preferred terms, non-preferred terms, related terms, etc.

Basically a thesaurus dictionary replaces all non-preferred terms by one preferred term and, optionally, preserves the original terms for indexing as well. PostgreSQL's current implementation of the thesaurus dictionary is an extension of the synonym dictionary with added phrase support. A thesaurus dictionary requires a configuration file of the following format:

```
# this is a comment
sample word(s) : indexed word(s)
mORE sAMPLE word(s) : more indexed word(s)
...
```

where the colon (:) symbol acts as a delimiter between a phrase and its replacement.

A thesaurus dictionary uses a subdictionary (which is specified in the dictionary's configuration) to normalize the input text before checking for phrase matches. It is only possible to select one subdictionary. An error is reported if the subdictionary fails to recognize a word. In that case, you should remove the use of the word or teach the subdictionary about it. You can place an asterisk (*) at the beginning of an indexed word to skip applying the subdictionary to it, but all sample words must be known to the subdictionary.

The thesaurus dictionary chooses the longest match if there are multiple phrases matching the input, and ties are broken by using the last definition.

Specific stop words recognized by the subdictionary cannot be specified; instead use ? to mark the location where any stop word can appear. For example, assuming that a and the are stop words:

```
? one ? two : swsw
```

matches a one the two and the one a two; both would be replaced by swsw.

Since a thesaurus dictionary has the capability to recognize phrases it must remember its state and interact with the parser. A thesaurus dictionary uses these assignments to check if it should handle the next word or stop accumulation. The thesaurus dictionary must be configured carefully. For example, if the thesaurus dictionary is assigned to handle only the asciiword token, then a thesaurus dictionary definition like one ? will not work since token type uint is not assigned to the thesaurus dictionary.

```
Caution
Thesauruses are used during indexing so any change in the thesaurus dictionary's parameters requires reindexing. For most other dictionary types, small changes such as adding or removing stopwords does not force reindexing.
```

12.6.4.1. Thesaurus Configuration

To define a new thesaurus dictionary, use the thesaurus template. For example:

```
CREATE TEXT SEARCH DICTIONARY thesaurus_simple (
FULL TEXT SEARCH

TEMPLATE = thesaurus,
DictFile = mythesaurus,
Dictionary = pg_catalog.english_stem
);

Here:

- thesaurus_simple is the new dictionary's name
- mythesaurus is the base name of the thesaurus configuration file. (Its full name will be $SHAREDIR/tsearch_data/mythesaurus.ths, where $SHAREDIR means the installation shared-data directory.)
- pg_catalog.english_stem is the subdictionary (here, a Snowball English stemmer) to use for thesaurus normalization. Notice that the subdictionary will have its own configuration (for example, stop words), which is not shown here.

Now it is possible to bind the thesaurus dictionary thesaurus_simple to the desired token types in a configuration, for example:

ALTER TEXT SEARCH CONFIGURATION russian
    ALTER MAPPING FOR asciword, asciihword, hword_asciipart
    WITH thesaurus_simple;

12.6.4.2. Thesaurus Example

Consider a simple astronomical thesaurus thesaurus_astro, which contains some astronomical word combinations:

supernovae stars : sn
crab nebulae : crab

Below we create a dictionary and bind some token types to an astronomical thesaurus and English stemmer:

CREATE TEXT SEARCH DICTIONARY thesaurus_astro (  
    TEMPLATE = thesaurus,
    DictFile = thesaurus_astro,
    Dictionary = english_stem
);

ALTER TEXT SEARCH CONFIGURATION russian
    ALTER MAPPING FOR asciword, asciihword, hword_asciipart
    WITH thesaurus_astro, english_stem;

Now we can see how it works. ts_lexize is not very useful for testing a thesaurus, because it treats its input as a single token. Instead we can use plainto_tsquery and to_tsvector which will break their input strings into multiple tokens:

SELECT plainto_tsquery('supernova star');
plainto_tsquery
-----------------
'sn'

SELECT to_tsvector('supernova star');
to_tsvector
----------
'sn':1

In principle, one can use to_tsquery if you quote the argument:

```
SELECT to_tsquery('''supernova star''');
to_tsquery
------------
'sn'
```

Notice that supernova star matches supernovae stars in thesaurus_astro because we specified the english_stem stemmer in the thesaurus definition. The stemmer removed the e and s.

To index the original phrase as well as the substitute, just include it in the right-hand part of the definition:

```
supernovae stars : sn supernovae stars
```

```
SELECT plainto_tsquery('supernova star');
plainto_tsquery
-------------------------------
'sn' & 'supernova' & 'star'
```

### 12.6.5. Ispell Dictionary

The Ispell dictionary template supports morphological dictionaries, which can normalize many different linguistic forms of a word into the same lexeme. For example, an English Ispell dictionary can match all declensions and conjugations of the search term bank, e.g., banking, banked, banks, banks’, and bank’s.

The standard PostgreSQL distribution does not include any Ispell configuration files. Dictionaries for a large number of languages are available from Ispell. Also, some more modern dictionary file formats are supported — MySpell (OO < 2.0.1) and Hunspell (OO >= 2.0.2). A large list of dictionaries is available on the OpenOffice Wiki.

To create an Ispell dictionary perform these steps:

- download dictionary configuration files. OpenOffice extension files have the .oxt extension. It is necessary to extract .aff and .dic files, change extensions to .affix and .dict. For some dictionary files it is also needed to convert characters to the UTF-8 encoding with commands (for example, for a Norwegian language dictionary):

  ```
  iconv -f ISO_8859-1 -t UTF-8 -o nn_no.affix nn_NO.aff
  iconv -f ISO_8859-1 -t UTF-8 -o nn_no.dict nn_NO.dic
  ```

- copy files to the $SHAREDIR/tsearch_data directory

- load files into PostgreSQL with the following command:

  ```
  CREATE TEXT SEARCH DICTIONARY english_hunspell ( 
    TEMPLATE = ispell, 
    DictFile = en_us, 
    AffFile = en_us, 
  )
  ```

---

1 https://www.cs.hmc.edu/~geoff/ispell.html
2 https://en.wikipedia.org/wiki/MySpell
3 https://sourceforge.net/projects/hunspell/
4 https://wiki.openoffice.org/wiki/Dictionaries
Stopwords = english);

Here, DictFile, AffFile, and StopWords specify the base names of the dictionary, affixes, and stop-words files. The stop-words file has the same format explained above for the simple dictionary type. The format of the other files is not specified here but is available from the above-mentioned web sites.

Ispell dictionaries usually recognize a limited set of words, so they should be followed by another broader dictionary; for example, a Snowball dictionary, which recognizes everything.

The .affix file of Ispell has the following structure:

prefixes
flag *A:
    .       > RE      # As in enter > reenter
suffixes
flag T:
    E       > ST      # As in late > latest
    [^AEIOU]Y > -Y,IEST # As in dirty > dirtiest
    [AEIOU]Y > EST     # As in gray > grayest
    [^EY]   > EST     # As in small > smallest

And the .dict file has the following structure:

lapse/ADGRS
lard/DGRS
large/PRTY
lark/MRS

Format of the .dict file is:

basic_form/affix_class_name

In the .affix file every affix flag is described in the following format:

condition > [-stripping_letters,] adding_affix

Here, condition has a format similar to the format of regular expressions. It can use groupings [...] and [^...]. For example, [AEIOU]Y means that the last letter of the word is "y" and the penultimate letter is "a", "e", "i", "o" or "u". [^EY] means that the last letter is neither "e" nor "y".

Ispell dictionaries support splitting compound words; a useful feature. Notice that the affix file should specify a special flag using the compounds controlled statement that marks dictionary words that can participate in compound formation:

compoundwords controlled z

Here are some examples for the Norwegian language:

SELECT ts_lexize('norwegian_ispell', 'overbuljongterningpakkmesterassistent');
{over,buljong,terning,pakk,mester,assistent}
SELECT ts_lexize('norwegian_ispell', 'sjokoladefabrikk');
{sjokoladefabrikk,sjokolade,fabrikk}
MySpell format is a subset of Hunspell. The `.affix` file of Hunspell has the following structure:

```
PFX A Y 1
PFX A 0 re .
SFX T N 4
SFX T 0 st e
SFX T y iest [^aeiou]y
SFX T 0 est [aeiou]y
SFX T 0 est [^ey]
```

The first line of an affix class is the header. Fields of an affix rules are listed after the header:

- parameter name (PFX or SFX)
- flag (name of the affix class)
- stripping characters from beginning (at prefix) or end (at suffix) of the word
- adding affix
- condition that has a format similar to the format of regular expressions.

The `.dict` file looks like the `.dict` file of Ispell:

```
larder/M
lardy/RT
large/RSPMYT
largehearted
```

**Note**

MySpell does not support compound words. Hunspell has sophisticated support for compound words. At present, PostgreSQL implements only the basic compound word operations of Hunspell.

### 12.6.6. Snowball Dictionary

The Snowball dictionary template is based on a project by Martin Porter, inventor of the popular Porter's stemming algorithm for the English language. Snowball now provides stemming algorithms for many languages (see the Snowball site\(^5\) for more information). Each algorithm understands how to reduce common variant forms of words to a base, or stem, spelling within its language. A Snowball dictionary requires a language parameter to identify which stemmer to use, and optionally can specify a stopword file name that gives a list of words to eliminate. (PostgreSQL’s standard stopword lists are also provided by the Snowball project.) For example, there is a built-in definition equivalent to

```sql
CREATE TEXT SEARCH DICTIONARY english_stem (    
    TEMPLATE = snowball, 
    Language = english, 
    StopWords = english 
) ;
```

The stopword file format is the same as already explained.

A Snowball dictionary recognizes everything, whether or not it is able to simplify the word, so it should be placed at the end of the dictionary list. It is useless to have it before any other dictionary because a token will never pass through it to the next dictionary.

---

\(^5\) [http://snowballstem.org/](http://snowballstem.org/)
12.7. Configuration Example

A text search configuration specifies all options necessary to transform a document into a tsvector: the parser to use to break text into tokens, and the dictionaries to use to transform each token into a lexeme. Every call of to_tsvector or to_tsquery needs a text search configuration to perform its processing. The configuration parameter default_text_search_config specifies the name of the default configuration, which is the one used by text search functions if an explicit configuration parameter is omitted. It can be set in postgresql.conf, or set for an individual session using the SET command.

Several predefined text search configurations are available, and you can create custom configurations easily. To facilitate management of text search objects, a set of SQL commands is available, and there are several psql commands that display information about text search objects (Section 12.10).

As an example we will create a configuration pg, starting by duplicating the built-in english configuration:

```
CREATE TEXT SEARCH CONFIGURATION public.pg ( COPY = pg_catalog.english );
```

We will use a PostgreSQL-specific synonym list and store it in $SHAREDIR/tsearch_data/pg_dict.syn. The file contents look like:

```
postgres    pg
pgsql       pg
postgresql  pg
```

We define the synonym dictionary like this:

```
CREATE TEXT SEARCH DICTIONARY pg_dict (
   TEMPLATE = synonym,
   SYNONYMS = pg_dict
);
```

Next we register the Ispell dictionary english_ispell, which has its own configuration files:

```
CREATE TEXT SEARCH DICTIONARY english_ispell (
   TEMPLATE = ispell,
   DictFile = english,
   AffFile = english,
   StopWords = english
);
```

Now we can set up the mappings for words in configuration pg:

```
ALTER TEXT SEARCH CONFIGURATION pg
   ALTER MAPPING FOR asciiword, asciihword, hword_asciipart,
               word, hword, hword_part
   WITH pg_dict, english_ispell, english_stem;
```

We choose not to index or search some token types that the built-in configuration does handle:

```
ALTER TEXT SEARCH CONFIGURATION pg
```
DROP MAPPING FOR email, url, url_path, sfloat, float;

Now we can test our configuration:

SELECT * FROM ts_debug('public.pg', 'PostgreSQL, the highly scalable, SQL compliant, open source object-relational database management system, is now undergoing beta testing of the next version of our software.');

The next step is to set the session to use the new configuration, which was created in the public schema:

=> \dF
  List of text search configurations
  Schema | Name | Description
          |------|-----------------|
  public  | pg   |                 |

SET default_text_search_config = 'public.pg';
SET SHOW default_text_search_config;
default_text_search_config
-----------
public.pg

12.8. Testing and Debugging Text Search

The behavior of a custom text search configuration can easily become confusing. The functions described in this section are useful for testing text search objects. You can test a complete configuration, or test parsers and dictionaries separately.

12.8.1. Configuration Testing

The function ts_debug allows easy testing of a text search configuration.

\text{ts\_debug][] config regconfig, ] document text,}
\hspace{1cm} OUT alias text,\hspace{1cm} \text{OUT description text,}
\hspace{1cm} OUT token text,\hspace{1cm} \text{OUT dictionaries regdictionary[],}
\hspace{1cm} OUT dictionary regdictionary,\hspace{1cm} \text{OUT lexemes text[]})
\hspace{1cm} \text{returns setof record}\n
\text{ts\_debug displays information about every token of document as produced by the parser and processed by the configured dictionaries. It uses the configuration specified by config, or default\_text\_search\_config if that argument is omitted.}
\n\text{ts\_debug returns one row for each token identified in the text by the parser. The columns returned are}
• **alias text** — short name of the token type
• **description text** — description of the token type
• **token text** — text of the token
• **dictionaries regdictionary[]** — the dictionaries selected by the configuration for this token type
• **dictionary regdictionary** — the dictionary that recognized the token, or NULL if none did
• **lexemes text[]** — the lexeme(s) produced by the dictionary that recognized the token, or NULL if none did; an empty array ({} ) means it was recognized as a stop word

Here is a simple example:

```sql
SELECT * FROM ts_debug('english','a fat cat sat on a mat - it ate a fat rats');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
<th>dictionaries</th>
<th>dictionary</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>a</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td>{}</td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>fat</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{fat}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>cat</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{cat}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>sat</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{sat}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>on</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>a</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>mat</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{mat}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td>-</td>
<td>{}</td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>it</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>ate</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{ate}</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>a</td>
<td>{english_stem}</td>
<td></td>
</tr>
<tr>
<td>english_stem</td>
<td>{}</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
For a more extensive demonstration, we first create a public.english configuration and Ispell dictionary for the English language:

```sql
CREATE TEXT SEARCH CONFIGURATION public.english ( COPY = pg_catalog.english );
CREATE TEXT SEARCH DICTIONARY english_ispell ( TEMPLATE = ispell, DictFile = english, AffFile = english, StopWords = english );
ALTER TEXT SEARCH CONFIGURATION public.english ALTER MAPPING FOR asciiword WITH english_ispell, english_stem;
```

```sql
SELECT * FROM ts_debug('public.english','The Brightest supernovaes');
```

```
+-------------------------------+----------------+-------------
<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
<th>dictionaries</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>The</td>
<td>{english_ispell,english_stem}</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td>()</td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>Brightest</td>
<td>{english_ispell,english_stem}</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td>()</td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>supernovaes</td>
<td>{english_ispell,english_stem}</td>
</tr>
</tbody>
</table>
```

In this example, the word **Brightest** was recognized by the parser as an ASCII word (alias `asciiword`). For this token type the dictionary list is `english_ispell` and `english_stem`. The word was recognized by `english_ispell`, which reduced it to the noun **bright**. The word `supernovaes` is unknown to the `english_ispell` dictionary so it was passed to the next dictionary, and, fortunately, was recognized (in fact, `english_stem` is a Snowball dictionary which recognizes everything; that is why it was placed at the end of the dictionary list).

The word **The** was recognized by the `english_ispell` dictionary as a stop word (Section 12.6.1) and will not be indexed. The spaces are discarded too, since the configuration provides no dictionaries at all for them.

You can reduce the width of the output by explicitly specifying which columns you want to see:

```sql
SELECT alias, token, dictionary, lexemes
```
FROM ts_debug('public.english','The Brightest supernovaes');

<table>
<thead>
<tr>
<th>alias</th>
<th>token</th>
<th>dictionary</th>
<th>lexemes</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>The</td>
<td>english_ispell</td>
<td>{}</td>
</tr>
<tr>
<td>blank</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Brightest</td>
<td>english_ispell</td>
<td>{bright}</td>
</tr>
<tr>
<td>blank</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>supernovaes</td>
<td>english_stem</td>
<td>{supernova}</td>
</tr>
</tbody>
</table>

### 12.8.2. Parser Testing

The following functions allow direct testing of a text search parser.

```sql
SELECT * FROM ts_parse('default', '123 - a number');
```

<table>
<thead>
<tr>
<th>tokid</th>
<th>token</th>
</tr>
</thead>
<tbody>
<tr>
<td>22</td>
<td>123</td>
</tr>
<tr>
<td>12</td>
<td>-</td>
</tr>
<tr>
<td>1</td>
<td>a</td>
</tr>
<tr>
<td>12</td>
<td>number</td>
</tr>
</tbody>
</table>

```sql
SELECT * FROM ts_token_type('default');
```

<table>
<thead>
<tr>
<th>tokid</th>
<th>alias</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>asciiword</td>
<td>Word, all ASCII</td>
</tr>
<tr>
<td>2</td>
<td>word</td>
<td>Word, all letters</td>
</tr>
<tr>
<td>3</td>
<td>numword</td>
<td>Word, letters and digits</td>
</tr>
<tr>
<td>4</td>
<td>email</td>
<td>Email address</td>
</tr>
</tbody>
</table>
The `ts_lexize` function facilitates dictionary testing.

\[
\text{ts_lexize}(\text{dict regdictionary}, \text{token text}) \text{ returns text[]} 
\]

`ts_lexize` returns an array of lexemes if the input `token` is known to the dictionary, or an empty array if the token is known to the dictionary but it is a stop word, or `NULL` if it is an unknown word.

Examples:

```sql
SELECT ts_lexize('english_stem', 'stars');
\text{ts_lexize}
\text{---------}
\{\text{star}\}

SELECT ts_lexize('english_stem', 'a');
\text{ts_lexize}
\text{---------}
\{\}
```

**Note**

The `ts_lexize` function expects a single `token`, not text. Here is a case where this can be confusing:

```sql
SELECT ts_lexize('thesaurus_astro','supernovae stars')
\text{is null;}
?column?
\text{---------}
\text{t}
```

The thesaurus dictionary `thesaurus_astro` does know the phrase `supernovae stars`, but `ts_lexize` fails since it does not parse the input text but treats it as a
Full Text Search

12.9. GIN and GiST Index Types

There are two kinds of indexes that can be used to speed up full text searches. Note that indexes are not mandatory for full text searching, but in cases where a column is searched on a regular basis, an index is usually desirable.

CREATE INDEX name ON table USING GIN (column);

Creates a GIN (Generalized Inverted Index)-based index. The column must be of tsvector type.

CREATE INDEX name ON table USING GIST (column);

Creates a GiST (Generalized Search Tree)-based index. The column can be of tsvector or tsquery type.

GIN indexes are the preferred text search index type. As inverted indexes, they contain an index entry for each word (lexeme), with a compressed list of matching locations. Multi-word searches can find the first match, then use the index to remove rows that are lacking additional words. GIN indexes store only the words (lexemes) of tsvector values, and not their weight labels. Thus a table row recheck is needed when using a query that involves weights.

A GiST index is lossy, meaning that the index might produce false matches, and it is necessary to check the actual table row to eliminate such false matches. (PostgreSQL does this automatically when needed.) GiST indexes are lossy because each document is represented in the index by a fixed-length signature. The signature is generated by hashing each word into a single bit in an n-bit string, with all these bits OR-ed together to produce an n-bit document signature. When two words hash to the same bit position there will be a false match. If all words in the query have matches (real or false) then the table row must be retrieved to see if the match is correct.

Lossiness causes performance degradation due to unnecessary fetches of table records that turn out to be false matches. Since random access to table records is slow, this limits the usefulness of GiST indexes. The likelihood of false matches depends on several factors, in particular the number of unique words, so using dictionaries to reduce this number is recommended.

Note that GIN index build time can often be improved by increasing maintenance_work_mem, while GiST index build time is not sensitive to that parameter.

Partitioning of big collections and the proper use of GIN and GiST indexes allows the implementation of very fast searches with online update. Partitioning can be done at the database level using table inheritance, or by distributing documents over servers and collecting external search results, e.g. via Foreign Data access. The latter is possible because ranking functions use only local information.

12.10. psql Support

Information about text search configuration objects can be obtained in psql using a set of commands:

\dF{d,p,t} [+ | PATTERN]
An optional + produces more details.

The optional parameter \texttt{PATTERN} can be the name of a text search object, optionally schema-qualified. If \texttt{PATTERN} is omitted then information about all visible objects will be displayed. \texttt{PATTERN} can be a regular expression and can provide \textit{separate} patterns for the schema and object names. The following examples illustrate this:

\begin{verbatim}
=> \texttt{\df*fulltext*}
     List of text search configurations
    Schema | Name        | Description
    --------+--------------+-------------
    public  | fulltext_cfg |

=> \texttt{\df*.fulltext*}
     List of text search configurations
    Schema | Name        | Description
    --------+--------------+-------------
    fulltext| fulltext_cfg |
    public  | fulltext_cfg |
\end{verbatim}

The available commands are:

\begin{verbatim}
\texttt{\df[+] \{PATTERN\}}
List text search configurations (add + for more detail).

=> \texttt{\df russia}
     List of text search configurations
    Schema | Name        | Description
    --------+--------------+-------------
    pg_catalog | russian | configuration for russian language

=> \texttt{\df+ russia}
Text search configuration "pg_catalog.russian"
Parser: "pg_catalog.default"
    Token      | Dictionaries
    -----------------+--------------
    asciihword      | english_stem
    asciword        | english_stem
    email           | simple
    file            | simple
    float           | simple
    host            | simple
    hword           | russian_stem
    hword_asciipart | english_stem
    hword_numpart   | simple
    hword_part      | russian_stem
    int             | simple
    numhword        | simple
    numword         | simple
    sfloat          | simple
    uint            | simple
    url             | simple
    url_path        | simple
    version         | simple
    word            | russian_stem
\end{verbatim}
List text search dictionaries (add + for more detail).

```
=> \dFd

```

List of text search dictionaries

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_catalog</td>
<td>danish_stem</td>
<td>snowball stemmer for danish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>dutch_stem</td>
<td>snowball stemmer for dutch language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>english_stem</td>
<td>snowball stemmer for english language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>finnish_stem</td>
<td>snowball stemmer for finnish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>french_stem</td>
<td>snowball stemmer for french language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>german_stem</td>
<td>snowball stemmer for german language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>hungarian_stem</td>
<td>snowball stemmer for hungarian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>italian_stem</td>
<td>snowball stemmer for italian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>norwegian_stem</td>
<td>snowball stemmer for norwegian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>portuguese_stem</td>
<td>snowball stemmer for portuguese language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>romanian_stem</td>
<td>snowball stemmer for romanian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>russian_stem</td>
<td>snowball stemmer for russian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>simple</td>
<td>simple dictionary: just lower case and</td>
</tr>
<tr>
<td></td>
<td></td>
<td>check for stopword</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>spanish_stem</td>
<td>snowball stemmer for spanish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>swedish_stem</td>
<td>snowball stemmer for swedish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>turkish_stem</td>
<td>snowball stemmer for turkish language</td>
</tr>
</tbody>
</table>

```
\dFp[+] [PATTERN]

List text search parsers (add + for more detail).

```
=> \dFp

```

List of text search parsers

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_catalog</td>
<td>default</td>
<td>default word parser</td>
</tr>
</tbody>
</table>

```
\dFp+ [PATTERN]

Text search parser "pg_catalog.default"

<table>
<thead>
<tr>
<th>Method</th>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Start parse</td>
<td>prsd_start</td>
<td></td>
</tr>
</tbody>
</table>
Get next token | prsd_nexttoken |
End parse       | prsd_end       |
Get headline    | prsd_headline  |
Get token types | prsd_lextype   |

Token types for parser "pg_catalog.default"

<table>
<thead>
<tr>
<th>Token name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciihword</td>
<td>Hyphenated word, all ASCII</td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
</tr>
<tr>
<td>email</td>
<td>Email address</td>
</tr>
<tr>
<td>entity</td>
<td>XML entity</td>
</tr>
<tr>
<td>file</td>
<td>File or path name</td>
</tr>
<tr>
<td>float</td>
<td>Decimal notation</td>
</tr>
<tr>
<td>host</td>
<td>Host</td>
</tr>
<tr>
<td>hword</td>
<td>Hyphenated word, all letters</td>
</tr>
<tr>
<td>hword_asciipart</td>
<td>Hyphenated word part, all ASCII</td>
</tr>
<tr>
<td>hword_numpart</td>
<td>Hyphenated word part, letters and digits</td>
</tr>
<tr>
<td>hword_part</td>
<td>Hyphenated word part, all letters</td>
</tr>
<tr>
<td>int</td>
<td>Signed integer</td>
</tr>
<tr>
<td>numhword</td>
<td>Hyphenated word, letters and digits</td>
</tr>
<tr>
<td>numword</td>
<td>Word, letters and digits</td>
</tr>
<tr>
<td>protocol</td>
<td>Protocol head</td>
</tr>
<tr>
<td>sfloat</td>
<td>Scientific notation</td>
</tr>
<tr>
<td>tag</td>
<td>XML tag</td>
</tr>
<tr>
<td>uint</td>
<td>Unsigned integer</td>
</tr>
<tr>
<td>url</td>
<td>URL</td>
</tr>
<tr>
<td>url_path</td>
<td>URL path</td>
</tr>
<tr>
<td>version</td>
<td>Version number</td>
</tr>
<tr>
<td>word</td>
<td>Word, all letters</td>
</tr>
</tbody>
</table>

(23 rows)

\dFt[+] [PATTERN]

List text search templates (add + for more detail).

=> \dFt

List of text search templates

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_catalog</td>
<td>ispell</td>
<td>ispell dictionary</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>simple</td>
<td>simple dictionary: just lower case and</td>
</tr>
<tr>
<td></td>
<td></td>
<td>check for stopword</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>snowball</td>
<td>snowball stemmer</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>synonym</td>
<td>synonym dictionary: replace word by</td>
</tr>
<tr>
<td></td>
<td></td>
<td>its synonym</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>thesaurus</td>
<td>thesaurus dictionary: phrase by phrase</td>
</tr>
<tr>
<td></td>
<td></td>
<td>substitution</td>
</tr>
</tbody>
</table>

12.11. Limitations

The current limitations of PostgreSQL’s text search features are:

- The length of each lexeme must be less than 2K bytes
• The length of a tsvector (lexemes + positions) must be less than 1 megabyte
• The number of lexemes must be less than $2^{64}$
• Position values in tsvector must be greater than 0 and no more than 16,383
• The match distance in a $<N>$ (FOLLOWED BY) tsquery operator cannot be more than 16,384
• No more than 256 positions per lexeme
• The number of nodes (lexemes + operators) in a tsquery must be less than 32,768

For comparison, the PostgreSQL 8.1 documentation contained 10,441 unique words, a total of 335,420 words, and the most frequent word "postgresql" was mentioned 6,127 times in 655 documents.

Another example — the PostgreSQL mailing list archives contained 910,989 unique words with 57,491,343 lexemes in 461,020 messages.
Chapter 13. Concurrency Control

This chapter describes the behavior of the PostgreSQL database system when two or more sessions try to access the same data at the same time. The goals in that situation are to allow efficient access for all sessions while maintaining strict data integrity. Every developer of database applications should be familiar with the topics covered in this chapter.

13.1. Introduction

PostgreSQL provides a rich set of tools for developers to manage concurrent access to data. Internally, data consistency is maintained by using a multiversion model (Multiversion Concurrency Control, MVCC). This means that each SQL statement sees a snapshot of data (a database version) as it was some time ago, regardless of the current state of the underlying data. This prevents statements from viewing inconsistent data produced by concurrent transactions performing updates on the same data rows, providing transaction isolation for each database session. MVCC, by eschewing the locking methodologies of traditional database systems, minimizes lock contention in order to allow for reasonable performance in multiuser environments.

The main advantage of using the MVCC model of concurrency control rather than locking is that in MVCC locks acquired for querying (reading) data do not conflict with locks acquired for writing data, and so reading never blocks writing and writing never blocks reading. PostgreSQL maintains this guarantee even when providing the strictest level of transaction isolation through the use of an innovative Serializable Snapshot Isolation (SSI) level.

Table- and row-level locking facilities are also available in PostgreSQL for applications which don't generally need full transaction isolation and prefer to explicitly manage particular points of conflict. However, proper use of MVCC will generally provide better performance than locks. In addition, application-defined advisory locks provide a mechanism for acquiring locks that are not tied to a single transaction.

13.2. Transaction Isolation

The SQL standard defines four levels of transaction isolation. The most strict is Serializable, which is defined by the standard in a paragraph which says that any concurrent execution of a set of Serializable transactions is guaranteed to produce the same effect as running them one at a time in some order. The other three levels are defined in terms of phenomena, resulting from interaction between concurrent transactions, which must not occur at each level. The standard notes that due to the definition of Serializable, none of these phenomena are possible at that level. (This is hardly surprising -- if the effect of the transactions must be consistent with having been run one at a time, how could you see any phenomena caused by interactions?)

The phenomena which are prohibited at various levels are:

dirty read

A transaction reads data written by a concurrent uncommitted transaction.

nonrepeateable read

A transaction re-reads data it has previously read and finds that data has been modified by another transaction (that committed since the initial read).

phantom read

A transaction re-executes a query returning a set of rows that satisfy a search condition and finds that the set of rows satisfying the condition has changed due to another recently-committed transaction.
Concurrency Control

serialization anomaly

The result of successfully committing a group of transactions is inconsistent with all possible orderings of running those transactions one at a time.

The SQL standard and PostgreSQL-implemented transaction isolation levels are described in Table 13.1.

Table 13.1. Transaction Isolation Levels

<table>
<thead>
<tr>
<th>Isolation Level</th>
<th>Dirty Read</th>
<th>Nonrepeatable Read</th>
<th>Phantom Read</th>
<th>Serialization Anomaly</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read uncommitted</td>
<td>Allowed, but not in PG</td>
<td>Possible</td>
<td>Possible</td>
<td>Possible</td>
</tr>
<tr>
<td>Read committed</td>
<td>Not possible</td>
<td>Possible</td>
<td>Possible</td>
<td>Possible</td>
</tr>
<tr>
<td>Repeatable read</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Allowed, but not in PG</td>
<td>Possible</td>
</tr>
<tr>
<td>Serializable</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Not possible</td>
</tr>
</tbody>
</table>

In PostgreSQL, you can request any of the four standard transaction isolation levels, but internally only three distinct isolation levels are implemented, i.e. PostgreSQL’s Read Uncommitted mode behaves like Read Committed. This is because it is the only sensible way to map the standard isolation levels to PostgreSQL’s multiversion concurrency control architecture.

The table also shows that PostgreSQL’s Repeatable Read implementation does not allow phantom reads. Stricter behavior is permitted by the SQL standard: the four isolation levels only define which phenomena must not happen, not which phenomena must happen. The behavior of the available isolation levels is detailed in the following subsections.

To set the transaction isolation level of a transaction, use the command SET TRANSACTION.

Important

Some PostgreSQL data types and functions have special rules regarding transactional behavior. In particular, changes made to a sequence (and therefore the counter of a column declared using serial) are immediately visible to all other transactions and are not rolled back if the transaction that made the changes aborts. See Section 9.16 and Section 8.1.4.

13.2.1. Read Committed Isolation Level

Read Committed is the default isolation level in PostgreSQL. When a transaction uses this isolation level, a SELECT query (without a FOR UPDATE/SHARE clause) sees only data committed before the query began; it never sees either uncommitted data or changes committed during query execution by concurrent transactions. In effect, a SELECT query sees a snapshot of the database as of the instant the query begins to run. However, SELECT does see the effects of previous updates executed within its own transaction, even though they are not yet committed. Also note that two successive SELECT commands can see different data, even though they are within a single transaction, if other transactions commit changes after the first SELECT starts and before the second SELECT starts.

UPDATE, DELETE, SELECT FOR UPDATE, and SELECT FOR SHARE commands behave the same as SELECT in terms of searching for target rows: they will only find target rows that were committed as of the command start time. However, such a target row might have already been updated (or deleted or locked) by another concurrent transaction by the time it is found. In this case, the would-be updater will wait for the first updating transaction to commit or roll back (if it is still in progress). If the first
updater rolls back, then its effects are negated and the second updater can proceed with updating the originally found row. If the first updater commits, the second updater will ignore the row if the first updater deleted it, otherwise it will attempt to apply its operation to the updated version of the row. The search condition of the command (the WHERE clause) is re-evaluated to see if the updated version of the row still matches the search condition. If so, the second updater proceeds with its operation using the updated version of the row. In the case of SELECT FOR UPDATE and SELECT FOR SHARE, this means it is the updated version of the row that is locked and returned to the client.

**INSERT** with an ON CONFLICT DO UPDATE clause behaves similarly. In Read Committed mode, each row proposed for insertion will either insert or update. Unless there are unrelated errors, one of those two outcomes is guaranteed. If a conflict originates in another transaction whose effects are not yet visible to the INSERT, the UPDATE clause will affect that row, even though possibly no version of that row is conventionally visible to the command.

**INSERT** with an ON CONFLICT DO NOTHING clause may have insertion not proceed for a row due to the outcome of another transaction whose effects are not visible to the INSERT snapshot. Again, this is only the case in Read Committed mode.

Because of the above rules, it is possible for an updating command to see an inconsistent snapshot: it can see the effects of concurrent updating commands on the same rows it is trying to update, but it does not see effects of those commands on other rows in the database. This behavior makes Read Committed mode unsuitable for commands that involve complex search conditions; however, it is just right for simpler cases. For example, consider updating bank balances with transactions like:

```
BEGIN;
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 12345;
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 7534;
COMMIT;
```

If two such transactions concurrently try to change the balance of account 12345, we clearly want the second transaction to start with the updated version of the account's row. Because each command is affecting only a predetermined row, letting it see the updated version of the row does not create any troublesome inconsistency.

More complex usage can produce undesirable results in Read Committed mode. For example, consider a **DELETE** command operating on data that is being both added and removed from its restriction criteria by another command, e.g., assume **website** is a two-row table with **website.hits** equaling 9 and 10:

```
BEGIN;
UPDATE website SET hits = hits + 1;
-- run from another session: DELETE FROM website WHERE hits = 10;
COMMIT;
```

The **DELETE** will have no effect even though there is a **website.hits = 10** row before and after the **UPDATE**. This occurs because the pre-update row value 9 is skipped, and when the **UPDATE** completes and **DELETE** obtains a lock, the new row value is no longer 10 but 11, which no longer matches the criteria.

Because Read Committed mode starts each command with a new snapshot that includes all transactions committed up to that instant, subsequent commands in the same transaction will see the effects of the committed concurrent transaction in any case. The point at issue above is whether or not a single command sees an absolutely consistent view of the database.

The partial transaction isolation provided by Read Committed mode is adequate for many applications, and this mode is fast and simple to use; however, it is not sufficient for all cases. Applications that
do complex queries and updates might require a more rigorously consistent view of the database than
Read Committed mode provides.

## 13.2.2. Repeatable Read Isolation Level

The **Repeatable Read** isolation level only sees data committed before the transaction began; it nev-
er sees either uncommitted data or changes committed during transaction execution by concurrent
transactions. (However, the query does see the effects of previous updates executed within its own
transaction, even though they are not yet committed.) This is a stronger guarantee than is required by
the SQL standard for this isolation level, and prevents all of the phenomena described in Table 13.1
except for serialization anomalies. As mentioned above, this is specifically allowed by the standard,
which only describes the *minimum* protections each isolation level must provide.

This level is different from Read Committed in that a query in a repeatable read transaction sees a
snapshot as of the start of the first non-transaction-control statement in the transaction, not as of the
start of the current statement within the transaction. Thus, successive **SELECT** commands within a
*single* transaction see the same data, i.e., they do not see changes made by other transactions that
committed after their own transaction started.

Applications using this level must be prepared to retry transactions due to serialization failures.

**UPDATE**, **DELETE**, **SELECT FOR UPDATE**, and **SELECT FOR SHARE** commands behave the same
as **SELECT** in terms of searching for target rows: they will only find target rows that were committed
as of the transaction start time. However, such a target row might have already been updated (or deleted
or locked) by another concurrent transaction by the time it is found. In this case, the repeatable read
transaction will wait for the first updating transaction to commit or roll back (if it is still in progress). If
the first updater rolls back, then its effects are negated and the repeatable read transaction can proceed
with updating the originally found row. But if the first updater commits (and actually updated or
deleted the row, not just locked it) then the repeatable read transaction will be rolled back with the
message

ERROR: could not serialize access due to concurrent update

because a repeatable read transaction cannot modify or lock rows changed by other transactions after
the repeatable read transaction began.

When an application receives this error message, it should abort the current transaction and retry the
whole transaction from the beginning. The second time through, the transaction will see the previous-
ly-committed change as part of its initial view of the database, so there is no logical conflict in using
the new version of the row as the starting point for the new transaction's update.

Note that only updating transactions might need to be retried; read-only transactions will never have
serialization conflicts.

The Repeatable Read mode provides a rigorous guarantee that each transaction sees a completely stable
view of the database. However, this view will not necessarily always be consistent with some serial
(one at a time) execution of concurrent transactions of the same level. For example, even a read only
transaction at this level may see a control record updated to show that a batch has been completed but
not see one of the detail records which is logically part of the batch because it read an earlier revision
of the control record. Attempts to enforce business rules by transactions running at this isolation level
are not likely to work correctly without careful use of explicit locks to block conflicting transactions.

---

**Note**

Prior to PostgreSQL version 9.1, a request for the Serializable transaction isolation
level provided exactly the same behavior described here. To retain the legacy Serial-
izable behavior, Repeatable Read should now be requested.
13.2.3. Serializable Isolation Level

The *Serializable* isolation level provides the strictest transaction isolation. This level emulates serial transaction execution for all committed transactions; as if transactions had been executed one after another, serially, rather than concurrently. However, like the Repeatable Read level, applications using this level must be prepared to retry transactions due to serialization failures. In fact, this isolation level works exactly the same as Repeatable Read except that it monitors for conditions which could make execution of a concurrent set of serializable transactions behave in a manner inconsistent with all possible serial (one at a time) executions of those transactions. This monitoring does not introduce any blocking beyond that present in repeatable read, but there is some overhead to the monitoring, and detection of the conditions which could cause a *serialization anomaly* will trigger a *serialization failure*.

As an example, consider a table `mytab`, initially containing:

<table>
<thead>
<tr>
<th>class</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>10</td>
</tr>
<tr>
<td>1</td>
<td>20</td>
</tr>
<tr>
<td>2</td>
<td>100</td>
</tr>
<tr>
<td>2</td>
<td>200</td>
</tr>
</tbody>
</table>

Suppose that serializable transaction A computes:

```sql
SELECT SUM(value) FROM mytab WHERE class = 1;
```

and then inserts the result (30) as the value in a new row with `class = 2`. Concurrently, serializable transaction B computes:

```sql
SELECT SUM(value) FROM mytab WHERE class = 2;
```

and obtains the result 300, which it inserts in a new row with `class = 1`. Then both transactions try to commit. If either transaction were running at the Repeatable Read isolation level, both would be allowed to commit; but since there is no serial order of execution consistent with the result, using Serializable transactions will allow one transaction to commit and will roll the other back with this message:

```
ERROR: could not serialize access due to read/write dependencies among transactions
```

This is because if A had executed before B, B would have computed the sum 330, not 300, and similarly the other order would have resulted in a different sum computed by A.

When relying on Serializable transactions to prevent anomalies, it is important that any data read from a permanent user table not be considered valid until the transaction which read it has successfully committed. This is true even for read-only transactions, except that data read within a *deferrable* read-only transaction is known to be valid as soon as it is read, because such a transaction waits until it can acquire a snapshot guaranteed to be free from such problems before starting to read any data. In all other cases applications must not depend on results read during a transaction that later aborted; instead, they should retry the transaction until it succeeds.

To guarantee true serializability PostgreSQL uses *predicate locking*, which means that it keeps locks which allow it to determine when a write would have had an impact on the result of a previous read from a concurrent transaction, had it run first. In PostgreSQL these locks do not cause any blocking and therefore can *not* play any part in causing a deadlock. They are used to identify and flag dependencies.
Concurrency Control

among concurrent Serializable transactions which in certain combinations can lead to serialization anomalies. In contrast, a Read Committed or Repeatable Read transaction which wants to ensure data consistency may need to take out a lock on an entire table, which could block other users attempting to use that table, or it may use SELECT FOR UPDATE or SELECT FOR SHARE which not only can block other transactions but cause disk access.

Predicate locks in PostgreSQL, like in most other database systems, are based on data actually accessed by a transaction. These will show up in the pg_locks system view with a mode of SIReadLock. The particular locks acquired during execution of a query will depend on the plan used by the query, and multiple finer-grained locks (e.g., tuple locks) may be combined into fewer coarser-grained locks (e.g., page locks) during the course of the transaction to prevent exhaustion of the memory used to track the locks. A READ ONLY transaction may be able to release its SIRead locks before completion, if it detects that no conflicts can still occur which could lead to a serialization anomaly. In fact, READ ONLY transactions will often be able to establish that fact at startup and avoid taking any predicate locks. If you explicitly request a SERIALIZABLE READ ONLY DEFERRABLE transaction, it will block until it can establish this fact. (This is the only case where Serializable transactions block but Repeatable Read transactions don't.) On the other hand, SIRead locks often need to be kept past transaction commit, until overlapping read write transactions complete.

Consistent use of Serializable transactions can simplify development. The guarantee that any set of successfully committed concurrent Serializable transactions will have the same effect as if they were run one at a time means that if you can demonstrate that a single transaction, as written, will do the right thing when run by itself, you can have confidence that it will do the right thing in any mix of Serializable transactions, even without any information about what those other transactions might do, or it will not successfully commit. It is important that an environment which uses this technique have a generalized way of handling serialization failures (which always return with a SQLSTATE value of '40001'), because it will be very hard to predict exactly which transactions might contribute to the read/write dependencies and need to be rolled back to prevent serialization anomalies. The monitoring of read/write dependencies has a cost, as does the restart of transactions which are terminated with a serialization failure, but balanced against the cost and blocking involved in use of explicit locks and SELECT FOR UPDATE or SELECT FOR SHARE, Serializable transactions are the best performance choice for some environments.

While PostgreSQL's Serializable transaction isolation level only allows concurrent transactions to commit if it can prove there is a serial order of execution that would produce the same effect, it doesn't always prevent errors from being raised that would not occur in true serial execution. In particular, it is possible to see unique constraint violations caused by conflicts with overlapping Serializable transactions even after explicitly checking that the key isn't present before attempting to insert it. This can be avoided by making sure that all Serializable transactions that insert potentially conflicting keys explicitly check if they can do so first. For example, imagine an application that asks the user for a new key and then checks that it doesn't exist already by trying to select it first, or generates a new key by selecting the maximum existing key and adding one. If some Serializable transactions insert new keys directly without following this protocol, unique constraints violations might be reported even in cases where they could not occur in a serial execution of the concurrent transactions.

For optimal performance when relying on Serializable transactions for concurrency control, these issues should be considered:

- Declare transactions as READ ONLY when possible.
- Control the number of active connections, using a connection pool if needed. This is always an important performance consideration, but it can be particularly important in a busy system using Serializable transactions.
- Don't put more into a single transaction than needed for integrity purposes.
- Don't leave connections dangling “idle in transaction” longer than necessary. The configuration parameter idle_in_transaction_session_timeout may be used to automatically disconnect lingering sessions.
Concurrency Control

- Eliminate explicit locks, `SELECT FOR UPDATE`, and `SELECT FOR SHARE` where no longer needed due to the protections automatically provided by Serializable transactions.

- When the system is forced to combine multiple page-level predicate locks into a single relation-level predicate lock because the predicate lock table is short of memory, an increase in the rate of serialization failures may occur. You can avoid this by increasing `max_pred_locks_per_transaction`, `max_pred_locks_per_relation`, and/or `max_pred_locks_per_page`.

- A sequential scan will always necessitate a relation-level predicate lock. This can result in an increased rate of serialization failures. It may be helpful to encourage the use of index scans by reducing `random_page_cost` and/or increasing `cpu_tuple_cost`. Be sure to weigh any decrease in transaction rollbacks and restarts against any overall change in query execution time.

### 13.3. Explicit Locking

PostgreSQL provides various lock modes to control concurrent access to data in tables. These modes can be used for application-controlled locking in situations where MVCC does not give the desired behavior. Also, most PostgreSQL commands automatically acquire locks of appropriate modes to ensure that referenced tables are not dropped or modified in incompatible ways while the command executes. (For example, `TRUNCATE` cannot safely be executed concurrently with other operations on the same table, so it obtains an exclusive lock on the table to enforce that.)

To examine a list of the currently outstanding locks in a database server, use the `pg_locks` system view. For more information on monitoring the status of the lock manager subsystem, refer to Chapter 28.

#### 13.3.1. Table-level Locks

The list below shows the available lock modes and the contexts in which they are used automatically by PostgreSQL. You can also acquire any of these locks explicitly with the command LOCK. Remember that all of these lock modes are table-level locks, even if the name contains the word “row”; the names of the lock modes are historical. To some extent the names reflect the typical usage of each lock mode — but the semantics are all the same. The only real difference between one lock mode and another is the set of lock modes with which each conflicts (see Table 13.2). Two transactions cannot hold locks of conflicting modes on the same table at the same time. (However, a transaction never conflicts with itself. For example, it might acquire ACCESS EXCLUSIVE lock and later acquire ACCESS SHARE lock on the same table.) Non-conflicting lock modes can be held concurrently by many transactions. Notice in particular that some lock modes are self-conflicting (for example, an ACCESS EXCLUSIVE lock cannot be held by more than one transaction at a time) while others are not self-conflicting (for example, an ACCESS SHARE lock can be held by multiple transactions).

**Table-level Lock Modes**

**ACCESS SHARE**

Conflicts with the ACCESS EXCLUSIVE lock mode only.

The `SELECT` command acquires a lock of this mode on referenced tables. In general, any query that only *reads* a table and does not modify it will acquire this lock mode.

**ROW SHARE**

Conflicts with the EXCLUSIVE and ACCESS EXCLUSIVE lock modes.

The `SELECT FOR UPDATE` and `SELECT FOR SHARE` commands acquire a lock of this mode on the target table(s) (in addition to `ACCESS SHARE` locks on any other tables that are referenced but not selected `FOR UPDATE/FOR SHARE`).
Concurrent Control

ROW EXCLUSIVE

Conflicts with the SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes.

The commands UPDATE, DELETE, and INSERT acquire this lock mode on the target table (in addition to ACCESS SHARE locks on any other referenced tables). In general, this lock mode will be acquired by any command that modifies data in a table.

SHARE UPDATE EXCLUSIVE

Conflicts with the SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent schema changes and VACUUM runs.

Acquired by VACUUM (without FULL), ANALYZE, CREATE INDEX CONCURRENTLY, CREATE STATISTICS and ALTER TABLE VALIDATE and other ALTER TABLE variants (for full details see ALTER TABLE).

SHARE

Conflicts with the ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent data changes.

Acquired by CREATE INDEX (without CONCURRENTLY).

SHARE ROW EXCLUSIVE

Conflicts with the ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent data changes, and is self-exclusive so that only one session can hold it at a time.

Acquired by CREATE COLLATION, CREATE TRIGGER, and many forms of ALTER TABLE (see ALTER TABLE).

EXCLUSIVE

Conflicts with the ROW SHARE, ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode allows only concurrent ACCESS SHARE locks, i.e., only reads from the table can proceed in parallel with a transaction holding this lock mode.

Acquired by REFRESH MATERIALIZED VIEW CONCURRENTLY.

ACCESS EXCLUSIVE

Conflicts with locks of all modes (ACCESS SHARE, ROW SHARE, ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE). This mode guarantees that the holder is the only transaction accessing the table in any way.

Acquired by the DROP TABLE, TRUNCATE, REINDEX, CLUSTER, VACUUM FULL, and REFRESH MATERIALIZED VIEW (without CONCURRENTLY) commands. Many forms of ALTER TABLE also acquire a lock at this level. This is also the default lock mode for LOCK TABLE statements that do not specify a mode explicitly.

Tip

Only an ACCESS EXCLUSIVE lock blocks a SELECT (without FOR UPDATE/S IARE) statement.
Once acquired, a lock is normally held till end of transaction. But if a lock is acquired after establishing a savepoint, the lock is released immediately if the savepoint is rolled back to. This is consistent with the principle that ROLLBACK cancels all effects of the commands since the savepoint. The same holds for locks acquired within a PL/pgSQL exception block: an error escape from the block releases locks acquired within it.

### Table 13.2. Conflicting Lock Modes

<table>
<thead>
<tr>
<th>Requested Lock Mode</th>
<th>Current Lock Mode</th>
<th>Row Share</th>
<th>Row Exclusive</th>
<th>Row Exclusive with Update</th>
<th>Share</th>
<th>Share Exclusive</th>
<th>Share Exclusive with Update</th>
<th>Access</th>
<th>Access Exclusive</th>
<th>Access Exclusive with Update</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACCESS SHARE</td>
<td>ACCESS SHARE</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ROW SHARE</td>
<td>ROW SHARE</td>
<td></td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ROW EXCLUSIVE</td>
<td>ROW EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SHARE UPDATE EXCLUSIVE</td>
<td>SHARE UPDATE EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SHARE EXCLUSIVE</td>
<td>SHARE EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SHARE ROW EXCLUSIVE</td>
<td>SHARE ROW EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EXCLUSIVE</td>
<td>EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ACCESS EXCLUSIVE</td>
<td>ACCESS EXCLUSIVE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

### 13.3.2. Row-level Locks

In addition to table-level locks, there are row-level locks, which are listed as below with the contexts in which they are used automatically by PostgreSQL. See Table 13.3 for a complete table of row-level lock conflicts. Note that a transaction can hold conflicting locks on the same row, even in different subtransactions; but other than that, two transactions can never hold conflicting locks on the same row. Row-level locks do not affect data querying; they block only writers and lockers to the same row.

#### Row-level Lock Modes

**FOR UPDATE**

FOR UPDATE causes the rows retrieved by the SELECT statement to be locked as though for update. This prevents them from being locked, modified or deleted by other transactions until the current transaction ends. That is, other transactions that attempt UPDATE, DELETE, SELECT FOR UPDATE, SELECT FOR NO KEY UPDATE, SELECT FOR SHARE or SELECT FOR KEY SHARE of these rows will be blocked until the current transaction ends; conversely, SELECT FOR UPDATE will wait for a concurrent transaction that has run any of those commands on the same row, and will then lock and return the updated row (or no row, if the row was deleted). Within
Concurrency Control

a REPEATABLE READ or SERIALIZABLE transaction, however, an error will be thrown if a row to be locked has changed since the transaction started. For further discussion see Section 13.4.

The FOR UPDATE lock mode is also acquired by any DELETE on a row, and also by an UPDATE that modifies the values on certain columns. Currently, the set of columns considered for the UPDATE case are those that have a unique index on them that can be used in a foreign key (so partial indexes and expression indexes are not considered), but this may change in the future.

FOR NO KEY UPDATE

Behaves similarly to FOR UPDATE, except that the lock acquired is weaker: this lock will not block SELECT FOR KEY SHARE commands that attempt to acquire a lock on the same rows. This lock mode is also acquired by any UPDATE that does not acquire a FOR UPDATE lock.

FOR SHARE

Behaves similarly to FOR NO KEY UPDATE, except that it acquires a shared lock rather than exclusive lock on each retrieved row. A shared lock blocks other transactions from performing UPDATE, DELETE, SELECT FOR UPDATE or SELECT FOR NO KEY UPDATE on these rows, but it does not prevent them from performing SELECT FOR SHARE or SELECT FOR KEY SHARE.

FOR KEY SHARE

Behaves similarly to FOR SHARE, except that the lock is weaker: SELECT FOR UPDATE is blocked, but not SELECT FOR NO KEY UPDATE. A key-shared lock blocks other transactions from performing DELETE or any UPDATE that changes the key values, but not other UPDATE, and neither does it prevent SELECT FOR NO KEY UPDATE, SELECT FOR SHARE, or SELECT FOR KEY SHARE.

PostgreSQL doesn't remember any information about modified rows in memory, so there is no limit on the number of rows locked at one time. However, locking a row might cause a disk write, e.g., SELECT FOR UPDATE modifies selected rows to mark them locked, and so will result in disk writes.

Table 13.3. Conflicting Row-level Locks

<table>
<thead>
<tr>
<th>Requested Mode</th>
<th>Lock Mode</th>
<th>FOR SHARE</th>
<th>KEY FOR SHARE</th>
<th>FOR NO UPDATE</th>
<th>KEY FOR UPDATE</th>
</tr>
</thead>
<tbody>
<tr>
<td>FOR KEY SHARE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FOR SHARE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FOR NO KEY UPDATE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FOR UPDATE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
</tr>
</tbody>
</table>

13.3.3. Page-level Locks

In addition to table and row locks, page-level share/exclusive locks are used to control read/write access to table pages in the shared buffer pool. These locks are released immediately after a row is fetched or updated. Application developers normally need not be concerned with page-level locks, but they are mentioned here for completeness.

13.3.4. Deadlocks

The use of explicit locking can increase the likelihood of deadlocks, wherein two (or more) transactions each hold locks that the other wants. For example, if transaction 1 acquires an exclusive lock on table A and then tries to acquire an exclusive lock on table B, while transaction 2 has already...
exclusive-locked table B and now wants an exclusive lock on table A, then neither one can proceed. PostgreSQL automatically detects deadlock situations and resolves them by aborting one of the transactions involved, allowing the other(s) to complete. (Exactly which transaction will be aborted is difficult to predict and should not be relied upon.)

Note that deadlocks can also occur as the result of row-level locks (and thus, they can occur even if explicit locking is not used). Consider the case in which two concurrent transactions modify a table. The first transaction executes:

```
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 11111;
```

This acquires a row-level lock on the row with the specified account number. Then, the second transaction executes:

```
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 22222;
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 11111;
```

The first UPDATE statement successfully acquires a row-level lock on the specified row, so it succeeds in updating that row. However, the second UPDATE statement finds that the row it is attempting to update has already been locked, so it waits for the transaction that acquired the lock to complete. Transaction two is now waiting on transaction one to complete before it continues execution. Now, transaction one executes:

```
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 22222;
```

Transaction one attempts to acquire a row-level lock on the specified row, but it cannot: transaction two already holds such a lock. So it waits for transaction two to complete. Thus, transaction one is blocked on transaction two, and transaction two is blocked on transaction one: a deadlock condition. PostgreSQL will detect this situation and abort one of the transactions.

The best defense against deadlocks is generally to avoid them by being certain that all applications using a database acquire locks on multiple objects in a consistent order. In the example above, if both transactions had updated the rows in the same order, no deadlock would have occurred. One should also ensure that the first lock acquired on an object in a transaction is the most restrictive mode that will be needed for that object. If it is not feasible to verify this in advance, then deadlocks can be handled on-the-fly by retrying transactions that abort due to deadlocks.

So long as no deadlock situation is detected, a transaction seeking either a table-level or row-level lock will wait indefinitely for conflicting locks to be released. This means it is a bad idea for applications to hold transactions open for long periods of time (e.g., while waiting for user input).

### 13.3.5. Advisory Locks

PostgreSQL provides a means for creating locks that have application-defined meanings. These are called **advisory locks**, because the system does not enforce their use — it is up to the application to use them correctly. Advisory locks can be useful for locking strategies that are an awkward fit for the MVCC model. For example, a common use of advisory locks is to emulate pessimistic locking strategies typical of so-called “flat file” data management systems. While a flag stored in a table could be used for the same purpose, advisory locks are faster, avoid table bloat, and are automatically cleaned up by the server at the end of the session.

There are two ways to acquire an advisory lock in PostgreSQL: at session level or at transaction level. Once acquired at session level, an advisory lock is held until explicitly released or the session ends.
Unlike standard lock requests, session-level advisory lock requests do not honor transaction semantics: a lock acquired during a transaction that is later rolled back will still be held following the rollback, and likewise an unlock is effective even if the calling transaction fails later. A lock can be acquired multiple times by its owning process; for each completed lock request there must be a corresponding unlock request before the lock is actually released. Transaction-level lock requests, on the other hand, behave more like regular lock requests: they are automatically released at the end of the transaction, and there is no explicit unlock operation. This behavior is often more convenient than the session-level behavior for short-term usage of an advisory lock. Session-level and transaction-level lock requests for the same advisory lock identifier will block each other in the expected way. If a session already holds a given advisory lock, additional requests by it will always succeed, even if other sessions are awaiting the lock; this statement is true regardless of whether the existing lock hold and new request are at session level or transaction level.

Like all locks in PostgreSQL, a complete list of advisory locks currently held by any session can be found in the `pg_locks` system view.

Both advisory locks and regular locks are stored in a shared memory pool whose size is defined by the configuration variables `max_locks_per_transaction` and `max_connections`. Care must be taken not to exhaust this memory or the server will be unable to grant any locks at all. This imposes an upper limit on the number of advisory locks grantable by the server, typically in the tens to hundreds of thousands depending on how the server is configured.

In certain cases using advisory locking methods, especially in queries involving explicit ordering and `LIMIT` clauses, care must be taken to control the locks acquired because of the order in which SQL expressions are evaluated. For example:

```sql
SELECT pg_advisory_lock(id) FROM foo WHERE id = 12345; -- ok
SELECT pg_advisory_lock(id) FROM foo WHERE id > 12345 LIMIT 100; -- danger!
SELECT pg_advisory_lock(q.id) FROM
  (SELECT id FROM foo WHERE id > 12345 LIMIT 100)
q; -- ok
```

In the above queries, the second form is dangerous because the `LIMIT` is not guaranteed to be applied before the locking function is executed. This might cause some locks to be acquired that the application was not expecting, and hence would fail to release (until it ends the session). From the point of view of the application, such locks would be dangling, although still viewable in `pg_locks`.

The functions provided to manipulate advisory locks are described in Section 9.26.10.

### 13.4. Data Consistency Checks at the Application Level

It is very difficult to enforce business rules regarding data integrity using Read Committed transactions because the view of the data is shifting with each statement, and even a single statement may not restrict itself to the statement's snapshot if a write conflict occurs.

While a Repeatable Read transaction has a stable view of the data throughout its execution, there is a subtle issue with using MVCC snapshots for data consistency checks, involving something known as read/write conflicts. If one transaction writes data and a concurrent transaction attempts to read the same data (whether before or after the write), it cannot see the work of the other transaction. The reader then appears to have executed first regardless of which started first or which committed first. If that is as far as it goes, there is no problem, but if the reader also writes data which is read by a concurrent transaction there is now a transaction which appears to have run before either of the previously mentioned transactions. If the transaction which appears to have executed last actually
Concurrency Control

commits first, it is very easy for a cycle to appear in a graph of the order of execution of the transactions. When such a cycle appears, integrity checks will not work correctly without some help.

As mentioned in Section 13.2.3, Serializable transactions are just Repeatable Read transactions which add nonblocking monitoring for dangerous patterns of read/write conflicts. When a pattern is detected which could cause a cycle in the apparent order of execution, one of the transactions involved is rolled back to break the cycle.

13.4.1. Enforcing Consistency With Serializable Transactions

If the Serializable transaction isolation level is used for all writes and for all reads which need a consistent view of the data, no other effort is required to ensure consistency. Software from other environments which is written to use serializable transactions to ensure consistency should “just work” in this regard in PostgreSQL.

When using this technique, it will avoid creating an unnecessary burden for application programmers if the application software goes through a framework which automatically retries transactions which are rolled back with a serialization failure. It may be a good idea to set `default_transaction_isolation` to `serializable`. It would also be wise to take some action to ensure that no other transaction isolation level is used, either inadvertently or to subvert integrity checks, through checks of the transaction isolation level in triggers.

See Section 13.2.3 for performance suggestions.

Warning

This level of integrity protection using Serializable transactions does not yet extend to hot standby mode (Section 26.5). Because of that, those using hot standby may want to use Repeatable Read and explicit locking on the master.

13.4.2. Enforcing Consistency With Explicit Blocking Locks

When non-serializable writes are possible, to ensure the current validity of a row and protect it against concurrent updates one must use `SELECT FOR UPDATE`, `SELECT FOR SHARE`, or an appropriate `LOCK TABLE` statement. (`SELECT FOR UPDATE` and `SELECT FOR SHARE` lock just the returned rows against concurrent updates, while `LOCK TABLE` locks the whole table.) This should be taken into account when porting applications to PostgreSQL from other environments.

Also of note to those converting from other environments is the fact that `SELECT FOR UPDATE` does not ensure that a concurrent transaction will not update or delete a selected row. To do that in PostgreSQL you must actually update the row, even if no values need to be changed. `SELECT FOR UPDATE temporarily blocks` other transactions from acquiring the same lock or executing an `UPDATE` or `DELETE` which would affect the locked row, but once the transaction holding this lock commits or rolls back, a blocked transaction will proceed with the conflicting operation unless an actual `UPDATE` of the row was performed while the lock was held.

Global validity checks require extra thought under non-serializable MVCC. For example, a banking application might wish to check that the sum of all credits in one table equals the sum of debits in another table, when both tables are being actively updated. Comparing the results of two successive `SELECT sum(...)` commands will not work reliably in Read Committed mode, since the second query will likely include the results of transactions not counted by the first. Doing the two sums in a single repeatable read transaction will give an accurate picture of only the effects of transactions that committed before the repeatable read transaction started — but one might legitimately wonder whether
Concurrency Control

the answer is still relevant by the time it is delivered. If the repeatable read transaction itself applied some changes before trying to make the consistency check, the usefulness of the check becomes even more debatable, since now it includes some but not all post-transaction-start changes. In such cases a careful person might wish to lock all tables needed for the check, in order to get an indisputable picture of current reality. A SHARE mode (or higher) lock guarantees that there are no uncommitted changes in the locked table, other than those of the current transaction.

Note also that if one is relying on explicit locking to prevent concurrent changes, one should either use Read Committed mode, or in Repeatable Read mode be careful to obtain locks before performing queries. A lock obtained by a repeatable read transaction guarantees that no other transactions modifying the table are still running, but if the snapshot seen by the transaction predates obtaining the lock, it might predate some now-committed changes in the table. A repeatable read transaction's snapshot is actually frozen at the start of its first query or data-modification command (SELECT, INSERT, UPDATE, or DELETE), so it is possible to obtain locks explicitly before the snapshot is frozen.

13.5. Caveats

Some DDL commands, currently only TRUNCATE and the table-rewriting forms of ALTER TABLE, are not MVCC-safe. This means that after the truncation or rewrite commits, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the DDL command committed. This will only be an issue for a transaction that did not access the table in question before the DDL command started — any transaction that has done so would hold at least an ACCESS SHARE table lock, which would block the DDL command until that transaction completes. So these commands will not cause any apparent inconsistency in the table contents for successive queries on the target table, but they could cause visible inconsistency between the contents of the target table and other tables in the database.

Support for the Serializable transaction isolation level has not yet been added to Hot Standby replication targets (described in Section 26.5). The strictest isolation level currently supported in hot standby mode is Repeatable Read. While performing all permanent database writes within Serializable transactions on the master will ensure that all standbys will eventually reach a consistent state, a Repeatable Read transaction run on the standby can sometimes see a transient state that is inconsistent with any serial execution of the transactions on the master.

13.6. Locking and Indexes

Though PostgreSQL provides nonblocking read/write access to table data, nonblocking read/write access is not currently offered for every index access method implemented in PostgreSQL. The various index types are handled as follows:

B-tree, GiST and SP-GiST indexes

Short-term share/exclusive page-level locks are used for read/write access. Locks are released immediately after each index row is fetched or inserted. These index types provide the highest concurrency without deadlock conditions.

Hash indexes

Share/exclusive hash-bucket-level locks are used for read/write access. Locks are released after the whole bucket is processed. Bucket-level locks provide better concurrency than index-level ones, but deadlock is possible since the locks are held longer than one index operation.

GIN indexes

Short-term share/exclusive page-level locks are used for read/write access. Locks are released immediately after each index row is fetched or inserted. But note that insertion of a GIN-indexed value usually produces several index key insertions per row, so GIN might do substantial work for a single value's insertion.
Currently, B-tree indexes offer the best performance for concurrent applications; since they also have more features than hash indexes, they are the recommended index type for concurrent applications that need to index scalar data. When dealing with non-scalar data, B-trees are not useful, and GiST, SP-GiST or GIN indexes should be used instead.
Chapter 14. Performance Tips

Query performance can be affected by many things. Some of these can be controlled by the user, while others are fundamental to the underlying design of the system. This chapter provides some hints about understanding and tuning PostgreSQL performance.

14.1. Using EXPLAIN

PostgreSQL devises a query plan for each query it receives. Choosing the right plan to match the query structure and the properties of the data is absolutely critical for good performance, so the system includes a complex planner that tries to choose good plans. You can use the EXPLAIN command to see what query plan the planner creates for any query. Plan-reading is an art that requires some experience to master, but this section attempts to cover the basics.

Examples in this section are drawn from the regression test database after doing a VACUUM ANALYZE, using 9.3 development sources. You should be able to get similar results if you try the examples yourself, but your estimated costs and row counts might vary slightly because ANALYZE’s statistics are random samples rather than exact, and because costs are inherently somewhat platform-dependent.

The examples use EXPLAIN’s default “text” output format, which is compact and convenient for humans to read. If you want to feed EXPLAIN’s output to a program for further analysis, you should use one of its machine-readable output formats (XML, JSON, or YAML) instead.

14.1.1. EXPLAIN Basics

The structure of a query plan is a tree of plan nodes. Nodes at the bottom level of the tree are scan nodes: they return raw rows from a table. There are different types of scan nodes for different table access methods: sequential scans, index scans, and bitmap index scans. There are also non-table row sources, such as VALUES clauses and set-returning functions in FROM, which have their own scan node types. If the query requires joining, aggregation, sorting, or other operations on the raw rows, then there will be additional nodes above the scan nodes to perform these operations. Again, there is usually more than one possible way to do these operations, so different node types can appear here too.

The output of EXPLAIN has one line for each node in the plan tree, showing the basic node type plus the cost estimates that the planner made for the execution of that plan node. Additional lines might appear, indented from the node's summary line, to show additional properties of the node. The very first line (the summary line for the topmost node) has the estimated total execution cost for the plan; it is this number that the planner seeks to minimize.

Here is a trivial example, just to show what the output looks like:

EXPLAIN SELECT * FROM tenk1;

QUERY PLAN
-------------
Seq Scan on tenk1  (cost=0.00..458.00 rows=10000 width=244)

Since this query has no WHERE clause, it must scan all the rows of the table, so the planner has chosen to use a simple sequential scan plan. The numbers that are quoted in parentheses are (left to right):

• Estimated start-up cost. This is the time expended before the output phase can begin, e.g., time to do the sorting in a sort node.

• Estimated total cost. This is stated on the assumption that the plan node is run to completion, i.e., all available rows are retrieved. In practice a node's parent node might stop short of reading all available rows (see the LIMIT example below).
• Estimated number of rows output by this plan node. Again, the node is assumed to be run to completion.

• Estimated average width of rows output by this plan node (in bytes).

The costs are measured in arbitrary units determined by the planner's cost parameters (see Section 19.7.2). Traditional practice is to measure the costs in units of disk page fetches; that is, seq_page_cost is conventionally set to 1.0 and the other cost parameters are set relative to that. The examples in this section are run with the default cost parameters.

It's important to understand that the cost of an upper-level node includes the cost of all its child nodes. It's also important to realize that the cost only reflects things that the planner cares about. In particular, the cost does not consider the time spent transmitting result rows to the client, which could be an important factor in the real elapsed time; but the planner ignores it because it cannot change it by altering the plan. (Every correct plan will output the same row set, we trust.)

The rows value is a little tricky because it is not the number of rows processed or scanned by the plan node, but rather the number emitted by the node. This is often less than the number scanned, as a result of filtering by any WHERE-clause conditions that are being applied at the node. Ideally the top-level rows estimate will approximate the number of rows actually returned, updated, or deleted by the query.

Returning to our example:

```
EXPLAIN SELECT * FROM tenk1;
```

```
QUERY PLAN
-------------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..458.00 rows=10000 width=244)
```

These numbers are derived very straightforwardly. If you do:

```
SELECT relpages, reltuples FROM pg_class WHERE relname = 'tenk1';
```

you will find that tenk1 has 358 disk pages and 10000 rows. The estimated cost is computed as (disk pages read * seq_page_cost) + (rows scanned * cpu_tuple_cost). By default, seq_page_cost is 1.0 and cpu_tuple_cost is 0.01, so the estimated cost is (358 * 1.0) + (10000 * 0.01) = 458.

Now let's modify the query to add a WHERE condition:

```
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 7000;
```

```
QUERY PLAN
------------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..483.00 rows=7001 width=244)
Filter: (unique1 < 7000)
```

Notice that the EXPLAIN output shows the WHERE clause being applied as a "filter" condition attached to the Seq Scan plan node. This means that the plan node checks the condition for each row it scans, and outputs only the ones that pass the condition. The estimate of output rows has been reduced because of the WHERE clause. However, the scan will still have to visit all 10000 rows, so the cost hasn't decreased; in fact it has gone up a bit (by 10000 * cpu_operator_cost, to be exact) to reflect the extra CPU time spent checking the WHERE condition.

The actual number of rows this query would select is 7000, but the rows estimate is only approximate. If you try to duplicate this experiment, you will probably get a slightly different estimate; moreover, it can change after each ANALYZE command, because the statistics produced by ANALYZE are taken from a randomized sample of the table.

Now, let's make the condition more restrictive:
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100;

 QUERY PLAN
------------------------------------------------------------------------------
Bitmap Heap Scan on tenk1  (cost=5.07..229.20 rows=101 width=244)  
  Recheck Cond: (unique1 < 100)  
  ->  Bitmap Index Scan on tenk1_unique1  (cost=0.00..5.04 rows=101 width=0)  
      Index Cond: (unique1 < 100)

Here the planner has decided to use a two-step plan: the child plan node visits an index to find the locations of rows matching the index condition, and then the upper plan node actually fetches those rows from the table itself. Fetching rows separately is much more expensive than reading them sequentially, but because not all the pages of the table have to be visited, this is still cheaper than a sequential scan. (The reason for using two plan levels is that the upper plan node sorts the row locations identified by the index into physical order before reading them, to minimize the cost of separate fetches. The “bitmap” mentioned in the node names is the mechanism that does the sorting.)

Now let's add another condition to the WHERE clause:

EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND stringu1 = 'xxx';

 QUERY PLAN
------------------------------------------------------------------------------
Bitmap Heap Scan on tenk1  (cost=5.04..229.43 rows=1 width=244)  
  Recheck Cond: (unique1 < 100)  
  Filter: (stringu1 = 'xxx'::name)  
  ->  Bitmap Index Scan on tenk1_unique1  (cost=0.00..5.04 rows=101 width=0)  
      Index Cond: (unique1 < 100)

The added condition stringu1 = 'xxx' reduces the output row count estimate, but not the cost because we still have to visit the same set of rows. Notice that the stringu1 clause cannot be applied as an index condition, since this index is only on the unique1 column. Instead it is applied as a filter on the rows retrieved by the index. Thus the cost has actually gone up slightly to reflect this extra checking.

In some cases the planner will prefer a “simple” index scan plan:

EXPLAIN SELECT * FROM tenk1 WHERE unique1 = 42;

 QUERY PLAN
------------------------------------------------------------------------------
Index Scan using tenk1_unique1 on tenk1  (cost=0.29..8.30 rows=1 width=244)  
  Index Cond: (unique1 = 42)

In this type of plan the table rows are fetched in index order, which makes them even more expensive to read, but there are so few that the extra cost of sorting the row locations is not worth it. You'll most often see this plan type for queries that fetch just a single row. It's also often used for queries that have an ORDER BY condition that matches the index order, because then no extra sorting step is needed to satisfy the ORDER BY.

If there are separate indexes on several of the columns referenced in WHERE, the planner might choose to use an AND or OR combination of the indexes:
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000;

QUERY PLAN

Bitmap Heap Scan on tenk1 (cost=25.08..60.21 rows=10 width=244)
  Recheck Cond: ((unique1 < 100) AND (unique2 > 9000))
  -> BitmapAnd (cost=25.08..25.08 rows=10 width=0)
    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
      Index Cond: (unique1 < 100)
    -> Bitmap Index Scan on tenk1_unique2 (cost=0.00..19.78 rows=999 width=0)
      Index Cond: (unique2 > 9000)

But this requires visiting both indexes, so it’s not necessarily a win compared to using just one index and treating the other condition as a filter. If you vary the ranges involved you’ll see the plan change accordingly.

Here is an example showing the effects of LIMIT:

EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000 LIMIT 2;

QUERY PLAN

Limit (cost=0.29..14.48 rows=2 width=244)
  -> Index Scan using tenk1_unique2 on tenk1 (cost=0.29..71.27 rows=10 width=244)
    Index Cond: (unique2 > 9000)
    Filter: (unique1 < 100)

This is the same query as above, but we added a LIMIT so that not all the rows need be retrieved, and the planner changed its mind about what to do. Notice that the total cost and row count of the Index Scan node are shown as if it were run to completion. However, the Limit node is expected to stop after retrieving only a fifth of those rows, so its total cost is only a fifth as much, and that’s the actual estimated cost of the query. This plan is preferred over adding a Limit node to the previous plan because the Limit could not avoid paying the startup cost of the bitmap scan, so the total cost would be something over 25 units with that approach.

Let’s try joining two tables, using the columns we have been discussing:

EXPLAIN SELECT *
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 10 AND t1.unique2 = t2.unique2;

QUERY PLAN

Nested Loop (cost=4.65..118.62 rows=10 width=488)
  -> Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244)
    Recheck Cond: (unique1 < 10)
    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0)
      Index Cond: (unique1 < 10)
  -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..7.91 rows=1 width=244)
Index Cond: (unique2 = t1.unique2)

In this plan, we have a nested-loop join node with two table scans as inputs, or children. The indentation of the node summary lines reflects the plan tree structure. The join's first, or “outer”, child is a bitmap scan similar to those we saw before. Its cost and row count are the same as we'd get from SELECT ... WHERE unique1 < 10 because we are applying the WHERE clause unique1 < 10 at that node. The t1.unique2 = t2.unique2 clause is not relevant yet, so it doesn't affect the row count of the outer scan. The nested-loop join node will run its second, or “inner” child once for each row obtained from the outer child. Column values from the current outer row can be plugged into the inner scan; here, the t1.unique2 value from the outer row is available, so we get a plan and costs similar to what we saw above for a simple SELECT ... WHERE t2.unique2 = constant case. (The estimated cost is actually a bit lower than what was seen above, as a result of caching that's expected to occur during the repeated index scans on t2.) The costs of the loop node are then set on the basis of the cost of the outer scan, plus one repetition of the inner scan for each outer row (10 * 7.91, here), plus a little CPU time for join processing.

In this example the join's output row count is the same as the product of the two scans' row counts, but that's not true in all cases because there can be additional WHERE clauses that mention both tables and so can only be applied at the join point, not to either input scan. Here's an example:

EXPLAIN SELECT *
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 10 AND t2.unique2 < 10 AND t1.hundred < t2.hundred;

QUERY PLAN

Nested Loop (cost=4.65..49.46 rows=33 width=488)
  Join Filter: (t1.hundred < t2.hundred)
    -> Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244)
      Recheck Cond: (unique1 < 10)
        -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0)
          Index Cond: (unique1 < 10)
    -> Materialize (cost=0.29..8.51 rows=10 width=244)
      -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..8.46 rows=10 width=244)
        Index Cond: (unique2 < 10)

The condition t1.hundred < t2.hundred can’t be tested in the tenk2_unique2 index, so it’s applied at the join node. This reduces the estimated output row count of the join node, but does not change either input scan.

Notice that here the planner has chosen to “materialize” the inner relation of the join, by putting a Materialize plan node atop it. This means that the t2 index scan will be done just once, even though the nested-loop join node needs to read that data ten times, once for each row from the outer relation. The Materialize node saves the data in memory as it’s read, and then returns the data from memory on each subsequent pass.

When dealing with outer joins, you might see join plan nodes with both “Join Filter” and plain “Filter” conditions attached. Join Filter conditions come from the outer join’s ON clause, so a row that fails the Join Filter condition could still get emitted as a null-extended row. But a plain Filter condition is applied after the outer-join rules and so acts to remove rows unconditionally. In an inner join there is no semantic difference between these types of filters.

If we change the query’s selectivity a bit, we might get a very different join plan:
EXPLAIN SELECT *
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

QUERY PLAN

Hash Join  (cost=230.47..713.98 rows=101 width=488)
 Hash Cond: (t2.unique2 = t1.unique2)
  ->  Seq Scan on tenk2 t2  (cost=0.00..445.00 rows=10000 width=244)
        Hash  (cost=229.20..229.20 rows=101 width=244)
                ->  Bitmap Heap Scan on tenk1 t1  (cost=5.07..229.20 rows=101 width=244)
                        Recheck Cond: (unique1 < 100)
                        ->  Bitmap Index Scan on tenk1_unique1  
                               (cost=0.00..5.04 rows=101 width=0)
                               Index Cond: (unique1 < 100)

Here, the planner has chosen to use a hash join, in which rows of one table are entered into an in-memory hash table, after which the other table is scanned and the hash table is probed for matches to each row. Again note how the indentation reflects the plan structure: the bitmap scan on tenk1 is the input to the Hash node, which constructs the hash table. That's then returned to the Hash Join node, which reads rows from its outer child plan and searches the hash table for each one.

Another possible type of join is a merge join, illustrated here:

EXPLAIN SELECT *
FROM tenk1 t1, onek t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

QUERY PLAN

Merge Join  (cost=198.11..268.19 rows=10 width=488)
 Merge Cond: (t1.unique2 = t2.unique2)
  ->  Index Scan using tenk1_unique2 on tenk1 t1  
       (cost=0.29..656.28 rows=101 width=244)
       Filter: (unique1 < 100)
  ->  Sort  (cost=197.83..200.33 rows=1000 width=244)
       Sort Key: t2.unique2
                ->  Seq Scan on onek t2  (cost=0.00..148.00 rows=1000 width=244)

Merge join requires its input data to be sorted on the join keys. In this plan the tenk1 data is sorted by using an index scan to visit the rows in the correct order, but a sequential scan and sort is preferred for onek, because there are many more rows to be visited in that table. (Sequential-scan-and-sort frequently beats an index scan for sorting many rows, because of the nonsequential disk access required by the index scan.)

One way to look at variant plans is to force the planner to disregard whatever strategy it thought was the cheapest, using the enable/disable flags described in Section 19.7.1. (This is a crude tool, but useful. See also Section 14.3.) For example, if we're unconvinced that sequential-scan-and-sort is the best way to deal with table onek in the previous example, we could try

SET enable_sort = off;

EXPLAIN SELECT *
FROM tenk1 t1, onek t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

### QUERY PLAN

```
Merge Join  (cost=0.56..292.65 rows=10 width=488)
  Merge Cond: (t1.unique2 = t2.unique2)
  ->  Index Scan using tenk1_unique2 on tenk1 t1  
      (cost=0.29..656.28 rows=101 width=244)
      Filter: (unique1 < 100)
  ->  Index Scan using onek_unique2 on onek t2  
      (cost=0.28..224.79 rows=1000 width=244)
```

which shows that the planner thinks that sorting `onek` by index-scanning is about 12% more expensive than sequential-scan-and-sort. Of course, the next question is whether it’s right about that. We can investigate that using `EXPLAIN ANALYZE`, as discussed below.

## 14.1.2. EXPLAIN ANALYZE

It is possible to check the accuracy of the planner’s estimates by using `EXPLAIN`’s `ANALYZE` option. With this option, `EXPLAIN` actually executes the query, and then displays the true row counts and true run time accumulated within each plan node, along with the same estimates that a plain `EXPLAIN` shows. For example, we might get a result like this:

```
EXPLAIN ANALYZE SELECT *
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 10 AND t1.unique2 = t2.unique2;

PLAN

```

Nested Loop  (cost=4.65..118.62 rows=10 width=488) (actual time=0.128..0.377 rows=10 loops=1)
  ->  Bitmap Heap Scan on tenk1 t1  
       (cost=4.36..39.47 rows=10 width=244) (actual time=0.057..0.121 rows=10 loops=1)
       Recheck Cond: (unique1 < 10)
  ->  Bitmap Index Scan on tenk1_unique1  
       (cost=0.00..4.36 rows=10 width=0) (actual time=0.024..0.024 rows=10 loops=1)
       Index Cond: (unique1 < 10)
  ->  Index Scan using tenk2_unique2 on tenk2 t2  
       (cost=0.29..7.91 rows=1000 width=244) (actual time=0.021..0.022 rows=1 loops=10)
       Index Cond: (unique2 = t1.unique2)
Planning time: 0.181 ms
Execution time: 0.501 ms
```

Note that the “actual time” values are in milliseconds of real time, whereas the cost estimates are expressed in arbitrary units; so they are unlikely to match up. The thing that’s usually most important to look for is whether the estimated row counts are reasonably close to reality. In this example the estimates were all dead-on, but that’s quite unusual in practice.

In some query plans, it is possible for a subplan node to be executed more than once. For example, the inner index scan will be executed once per outer row in the above nested-loop plan. In such cases, the `loops` value reports the total number of executions of the node, and the actual time and rows values shown are averages per-execution. This is done to make the numbers comparable with the way that the cost estimates are shown. Multiply by the `loops` value to get the total time actually spent in the node. In the above example, we spent a total of 0.220 milliseconds executing the index scans on `tenk2`.

In some cases `EXPLAIN ANALYZE` shows additional execution statistics beyond the plan node execution times and row counts. For example, Sort and Hash nodes provide extra information:
EXPLAIN ANALYZE SELECT * FROM tenk1 t1, tenk2 t2 WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2 ORDER BY t1.fivethous;

QUERY PLAN

Sort (cost=717.34..717.59 rows=101 width=488) (actual time=7.761..7.774 rows=100 loops=1)
  Sort Key: t1.fivethous
  Sort Method: quicksort  Memory: 77kB
  ->  Hash Join  (cost=230.47..713.98 rows=101 width=488) (actual time=0.711..7.427 rows=100 loops=1)
     Hash Cond: (t2.unique2 = t1.unique2)
     ->  Seq Scan on tenk2 t2  (cost=0.00..445.00 rows=10000 width=244) (actual time=0.007..2.583 rows=10000 loops=1)
     ->  Hash  (cost=229.20..229.20 rows=101 width=244) (actual time=0.659..0.659 rows=100 loops=1)
        Buckets: 1024  Batches: 1  Memory Usage: 28kB
        ->  Bitmap Heap Scan on tenk1 t1  (cost=5.07..229.20 rows=101 width=244) (actual time=0.080..0.526 rows=100 loops=1)
           Recheck Cond: (unique1 < 100)
           ->  Bitmap Index Scan on tenk1_unique1  (cost=0.00..5.04 rows=101 width=0) (actual time=0.049..0.049 rows=100 loops=1)
              Index Cond: (unique1 < 100)
Planning time: 0.194 ms
Execution time: 8.008 ms

The Sort node shows the sort method used (in particular, whether the sort was in-memory or on-disk) and the amount of memory or disk space needed. The Hash node shows the number of hash buckets and batches as well as the peak amount of memory used for the hash table. (If the number of batches exceeds one, there will also be disk space usage involved, but that is not shown.)

Another type of extra information is the number of rows removed by a filter condition:

EXPLAIN ANALYZE SELECT * FROM tenk1 WHERE ten < 7;

QUERY PLAN

Seq Scan on tenk1  (cost=0.00..483.00 rows=7000 width=244) (actual time=0.016..5.107 rows=7000 loops=1)
  Filter: (ten < 7)
  Rows Removed by Filter: 3000
Planning time: 0.083 ms
Execution time: 5.905 ms

These counts can be particularly valuable for filter conditions applied at join nodes. The “Rows Removed” line only appears when at least one scanned row, or potential join pair in the case of a join node, is rejected by the filter condition.

A case similar to filter conditions occurs with “lossy” index scans. For example, consider this search for polygons containing a specific point:

EXPLAIN ANALYZE SELECT * FROM polygon_tbl WHERE f1 @> polygon '(0.5,2.0)';
The planner thinks (quite correctly) that this sample table is too small to bother with an index scan, so we have a plain sequential scan in which all the rows got rejected by the filter condition. But if we force an index scan to be used, we see:

SET enable_seqscan TO off;

EXPLAIN ANALYZE SELECT * FROM polygon_tbl WHERE f1 @> polygon 
'((0.5,2.0)');

Here we can see that the index returned one candidate row, which was then rejected by a recheck of the index condition. This happens because a GiST index is "lossy" for polygon containment tests: it actually returns the rows with polygons that overlap the target, and then we have to do the exact containment test on those rows.

EXPLAIN has a BUFFERS option that can be used with ANALYZE to get even more run time statistics:

EXPLAIN (ANALYZE, BUFFERS) SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000;

Here is the plan for that query:

**QUERY PLAN**

Bitmap Heap Scan on tenk1  
(cost=25.08..60.21 rows=10 width=244) (actual time=0.323..0.342 rows=10 loops=1) 
  Recheck Cond: ((unique1 < 100) AND (unique2 > 9000)) 
  Buffers: shared hit=15 
  ->  BitmapAnd  
  (cost=25.08..25.08 rows=10 width=0) (actual time=0.309..0.309 rows=0 loops=1) 
  Buffers: shared hit=7 
    ->  Bitmap Index Scan on tenk1_unique1  
    (cost=0.00..5.04 rows=101 width=0) (actual time=0.043..0.043 rows=100 loops=1) 
    Index Cond: (unique1 < 100) 
    Buffers: shared hit=2 
      ->  Bitmap Index Scan on tenk1_unique2  
      (cost=0.00..19.78 rows=999 width=0) (actual time=0.227..0.227 rows=999 loops=1) 
      Index Cond: (unique2 > 9000) 
      Buffers: shared hit=5 
Planning time: 0.088 ms
Performance Tips

Execution time: 0.423 ms

The numbers provided by BUFFERS help to identify which parts of the query are the most I/O-intensive.

Keep in mind that because EXPLAIN ANALYZE actually runs the query, any side-effects will happen as usual, even though whatever results the query might output are discarded in favor of printing the EXPLAIN data. If you want to analyze a data-modifying query without changing your tables, you can roll the command back afterwards, for example:

```
BEGIN;
EXPLAIN ANALYZE UPDATE tenk1 SET hundred = hundred + 1 WHERE unique1 < 100;
```

As seen in this example, when the query is an INSERT, UPDATE, or DELETE command, the actual work of applying the table changes is done by a top-level Insert, Update, or Delete plan node. The plan nodes underneath this node perform the work of locating the old rows and/or computing the new data. So above, we see the same sort of bitmap table scan we’ve seen already, and its output is fed to an Update node that stores the updated rows. It’s worth noting that although the data-modifying node can take a considerable amount of run time (here, it’s consuming the lion’s share of the time), the planner does not currently add anything to the cost estimates to account for that work. That’s because the work to be done is the same for every correct query plan, so it doesn’t affect planning decisions.

When an UPDATE or DELETE command affects an inheritance hierarchy, the output might look like this:

```
EXPLAIN UPDATE parent SET f2 = f2 + 1 WHERE f1 = 101;
```

As seen in this example, when the query is an INSERT, UPDATE, or DELETE command, the actual work of applying the table changes is done by a top-level Insert, Update, or Delete plan node. The plan nodes underneath this node perform the work of locating the old rows and/or computing the new data. So above, we see the same sort of bitmap table scan we’ve seen already, and its output is fed to an Update node that stores the updated rows. It’s worth noting that although the data-modifying node can take a considerable amount of run time (here, it’s consuming the lion’s share of the time), the planner does not currently add anything to the cost estimates to account for that work. That’s because the work to be done is the same for every correct query plan, so it doesn’t affect planning decisions.
Index Cond: (f1 = 101)
  ->  Index Scan using child3_f1_key on child3  (cost=0.15..8.17 rows=1 width=14)
      Index Cond: (f1 = 101)
In this example the Update node needs to consider three child tables as well as the originally-mentioned parent table. So there are four input scanning subplans, one per table. For clarity, the Update node is annotated to show the specific target tables that will be updated, in the same order as the corresponding subplans. (These annotations are new as of PostgreSQL 9.5; in prior versions the reader had to intuit the target tables by inspecting the subplans.)

The Planning time shown by EXPLAIN ANALYZE is the time it took to generate the query plan from the parsed query and optimize it. It does not include parsing or rewriting.

The Execution time shown by EXPLAIN ANALYZE includes executor start-up and shut-down time, as well as the time to run any triggers that are fired, but it does not include parsing, rewriting, or planning time. Time spent executing BEFORE triggers, if any, is included in the time for the related Insert, Update, or Delete node; but time spent executing AFTER triggers is not counted there because AFTER triggers are fired after completion of the whole plan. The total time spent in each trigger (either BEFORE or AFTER) is also shown separately. Note that deferred constraint triggers will not be executed until end of transaction and are thus not considered at all by EXPLAIN ANALYZE.

14.1.3. Caveats

There are two significant ways in which run times measured by EXPLAIN ANALYZE can deviate from normal execution of the same query. First, since no output rows are delivered to the client, network transmission costs and I/O conversion costs are not included. Second, the measurement overhead added by EXPLAIN ANALYZE can be significant, especially on machines with slow gettimeofday() operating-system calls. You can use the pg_test_timing tool to measure the overhead of timing on your system.

EXPLAIN results should not be extrapolated to situations much different from the one you are actually testing; for example, results on a toy-sized table cannot be assumed to apply to large tables. The planner's cost estimates are not linear and so it might choose a different plan for a larger or smaller table. An extreme example is that on a table that only occupies one disk page, you'll nearly always get a sequential scan plan whether indexes are available or not. The planner realizes that it's going to take one disk page read to process the table in any case, so there's no value in expending additional page reads to look at an index. (We saw this happening in the polygon_tbl example above.)

There are cases in which the actual and estimated values won't match up well, but nothing is really wrong. One such case occurs when plan node execution is stopped short by a LIMIT or similar effect. For example, in the LIMIT query we used before,

EXPLAIN ANALYZE SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000 LIMIT 2;

QUERY

PLAN

Limit (cost=0.29..14.71 rows=2 width=244) (actual time=0.177..0.249 rows=2 loops=1)
  ->  Index Scan using tenk1_unique2 on tenk1  (cost=0.29..72.42 rows=10 width=244) (actual time=0.174..0.244 rows=2 loops=1)
      Index Cond: (unique2 > 9000)
      Filter: (unique1 < 100)
      Rows Removed by Filter: 287
Planning time: 0.096 ms
Execution time: 0.336 ms
the estimated cost and row count for the Index Scan node are shown as though it were run to comple-
tion. But in reality the Limit node stopped requesting rows after it got two, so the actual row count is
only 2 and the run time is less than the cost estimate would suggest. This is not an estimation error,
only a discrepancy in the way the estimates and true values are displayed.

Merge joins also have measurement artifacts that can confuse the unwary. A merge join will stop
reading one input if it’s exhausted the other input and the next key value in the one input is greater
than the last key value of the other input; in such a case there can be no more matches and so no need
to scan the rest of the first input. This results in not reading all of one child, with results like those
mentioned for LIMIT. Also, if the outer (first) child contains rows with duplicate key values, the
inner (second) child is backed up and rescanned for the portion of its rows matching that key value.
EXPLAIN ANALYZE counts these repeated emissions of the same inner rows as if they were real
additional rows. When there are many outer duplicates, the reported actual row count for the inner child
plan node can be significantly larger than the number of rows that are actually in the inner relation.

BitmapAnd and BitmapOr nodes always report their actual row counts as zero, due to implementa-
tion limitations.

Generally, the EXPLAIN output will display details for every plan node which was generated by the
query planner. However, there are cases where the executor is able to determine that certain nodes
are not required; currently, the only node type to support this is the Append node. This node type
has the ability to discard subnodes which it is able to determine won’t contain any records required by
the query. It is possible to determine that nodes have been removed in this way by the presence of a
"Subplans Removed" property in the EXPLAIN output.

14.2. Statistics Used by the Planner

14.2.1. Single-Column Statistics

As we saw in the previous section, the query planner needs to estimate the number of rows retrieved
by a query in order to make good choices of query plans. This section provides a quick look at the
statistics that the system uses for these estimates.

One component of the statistics is the total number of entries in each table and index, as well as
the number of disk blocks occupied by each table and index. This information is kept in the table
pg_class, in the columns reltuples and relpages. We can look at it with queries similar to
this one:

```
SELECT relname, relkind, reltuples, relpages
FROM pg_class
WHERE relname LIKE 'tenk1%';
```

```
<table>
<thead>
<tr>
<th>relname</th>
<th>relkind</th>
<th>reltuples</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>tenk1</td>
<td>r</td>
<td>10000</td>
<td>358</td>
</tr>
<tr>
<td>tenk1_hundred</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_thous_tenthous</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_unique1</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_unique2</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
</tbody>
</table>
```

Here we can see that tenk1 contains 10000 rows, as do its indexes, but the indexes are (unsurpris-
ingly) much smaller than the table.

For efficiency reasons, reltuples and relpages are not updated on-the-fly, and so they usually
contain somewhat out-of-date values. They are updated by VACUUM, ANALYZE, and a few DDL
commands such as CREATE INDEX. A VACUUM or ANALYZE operation that does not scan the entire
table (which is commonly the case) will incrementally update the reltuples count on the basis of the part of the table it did scan, resulting in an approximate value. In any case, the planner will scale the values it finds in pg_class to match the current physical table size, thus obtaining a closer approximation.

Most queries retrieve only a fraction of the rows in a table, due to WHERE clauses that restrict the rows to be examined. The planner thus needs to make an estimate of the selectivity of WHERE clauses, that is, the fraction of rows that match each condition in the WHERE clause. The information used for this task is stored in the pg_statistic system catalog. Entries in pg_statistic are updated by the ANALYZE and VACUUM ANALYZE commands, and are always approximate even when freshly updated.

Rather than look at pg_statistic directly, it's better to look at its view pg_stats when examining the statistics manually. pg_stats is designed to be more easily readable. Furthermore, pg_stats is readable by all, whereas pg_statistic is only readable by a superuser. (This prevents unprivileged users from learning something about the contents of other people's tables from the statistics. The pg_stats view is restricted to show only rows about tables that the current user can read.) For example, we might do:

```
SELECT attname, inherited, n_distinct,
     array_to_string(most_common_vals, E'\n') as most_common_vals
FROM pg_stats
WHERE tablename = 'road';
```

<table>
<thead>
<tr>
<th>attname</th>
<th>inherited</th>
<th>n_distinct</th>
<th>most_common_vals</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>f</td>
<td>-0.363388</td>
<td>I- 580</td>
</tr>
<tr>
<td>Ramp+</td>
<td></td>
<td></td>
<td>I- 880</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Sp Railroad</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I- 580</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I- 680</td>
</tr>
<tr>
<td>name</td>
<td>t</td>
<td>-0.284859</td>
<td>I- 880</td>
</tr>
<tr>
<td>Ramp+</td>
<td></td>
<td></td>
<td>I- 580</td>
</tr>
<tr>
<td>Ramp+</td>
<td></td>
<td></td>
<td>I- 580</td>
</tr>
<tr>
<td>Ramp+</td>
<td></td>
<td></td>
<td>I- 580</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>State Hwy 13</td>
</tr>
</tbody>
</table>

Note that two rows are displayed for the same column, one corresponding to the complete inheritance hierarchy starting at the road table (inherited=t), and another one including only the road table itself (inherited=f).

The amount of information stored in pg_statistic by ANALYZE, in particular the maximum number of entries in the most_common_vals and histogram_bounds arrays for each column, can be set on a column-by-column basis using the ALTER TABLE SET STATISTICS command.
or globally by setting the `default_statistics_target` configuration variable. The default limit is presently 100 entries. Raising the limit might allow more accurate planner estimates to be made, particularly for columns with irregular data distributions, at the price of consuming more space in `pg_statistic` and slightly more time to compute the estimates. Conversely, a lower limit might be sufficient for columns with simple data distributions.

Further details about the planner's use of statistics can be found in Chapter 70.

### 14.2.2. Extended Statistics

It is common to see slow queries running bad execution plans because multiple columns used in the query clauses are correlated. The planner normally assumes that multiple conditions are independent of each other, an assumption that does not hold when column values are correlated. Regular statistics, because of their per-individual-column nature, cannot capture any knowledge about cross-column correlation. However, PostgreSQL has the ability to compute *multivariate statistics*, which can capture such information.

Because the number of possible column combinations is very large, it's impractical to compute multivariate statistics automatically. Instead, *extended statistics objects*, more often called just *statistics objects*, can be created to instruct the server to obtain statistics across interesting sets of columns.

Statistics objects are created using `CREATE STATISTICS`, which see for more details. Creation of such an object merely creates a catalog entry expressing interest in the statistics. Actual data collection is performed by `ANALYZE` (either a manual command, or background auto-analyze). The collected values can be examined in the `pg_statistic_ext` catalog.

`ANALYZE` computes extended statistics based on the same sample of table rows that it takes for computing regular single-column statistics. Since the sample size is increased by increasing the statistics target for the table or any of its columns (as described in the previous section), a larger statistics target will normally result in more accurate extended statistics, as well as more time spent calculating them.

The following subsections describe the kinds of extended statistics that are currently supported.

#### 14.2.2.1. Functional Dependencies

The simplest kind of extended statistics tracks *functional dependencies*, a concept used in definitions of database normal forms. We say that column \( b \) is functionally dependent on column \( a \) if knowledge of the value of \( a \) is sufficient to determine the value of \( b \), that is there are no two rows having the same value of \( a \) but different values of \( b \). In a fully normalized database, functional dependencies should exist only on primary keys and superkeys. However, in practice many data sets are not fully normalized for various reasons; intentional denormalization for performance reasons is a common example. Even in a fully normalized database, there may be partial correlation between some columns, which can be expressed as partial functional dependency.

The existence of functional dependencies directly affects the accuracy of estimates in certain queries. If a query contains conditions on both the independent and the dependent column(s), the conditions on the dependent columns do not further reduce the result size; but without knowledge of the functional dependency, the query planner will assume that the conditions are independent, resulting in underestimating the result size.

To inform the planner about functional dependencies, `ANALYZE` can collect measurements of cross-column dependency. Assessing the degree of dependency between all sets of columns would be prohibitively expensive, so data collection is limited to those groups of columns appearing together in a statistics object defined with the `dependencies` option. It is advisable to create `dependencies` statistics only for column groups that are strongly correlated, to avoid unnecessary overhead in both `ANALYZE` and later query planning.

Here is an example of collecting functional-dependency statistics:
CREATE STATISTICS stts (dependencies) ON zip, city FROM zipcodes;

ANALYZE zipcodes;

SELECT stxname, stxkeys, stxdependencies
FROM pg_statistic_ext
WHERE stxname = 'stts';

<table>
<thead>
<tr>
<th>stxname</th>
<th>stxkeys</th>
<th>stxdependencies</th>
</tr>
</thead>
<tbody>
<tr>
<td>stts</td>
<td>1 5</td>
<td>{&quot;1 =&gt; 5&quot;: 1.000000, &quot;5 =&gt; 1&quot;: 0.423130}</td>
</tr>
</tbody>
</table>

(1 row)

Here it can be seen that column 1 (zip code) fully determines column 5 (city) so the coefficient is 1.0, while city only determines zip code about 42% of the time, meaning that there are many cities (58%) that are represented by more than a single ZIP code.

When computing the selectivity for a query involving functionally dependent columns, the planner adjusts the per-condition selectivity estimates using the dependency coefficients so as not to produce an underestimate.

### 14.2.2.1.1. Limitations of Functional Dependencies

Functional dependencies are currently only applied when considering simple equality conditions that compare columns to constant values. They are not used to improve estimates for equality conditions comparing two columns or comparing a column to an expression, nor for range clauses, LIKE or any other type of condition.

When estimating with functional dependencies, the planner assumes that conditions on the involved columns are compatible and hence redundant. If they are incompatible, the correct estimate would be zero rows, but that possibility is not considered. For example, given a query like

```
SELECT * FROM zipcodes WHERE city = 'San Francisco' AND zip = '94105';
```

the planner will disregard the city clause as not changing the selectivity, which is correct. However, it will make the same assumption about

```
SELECT * FROM zipcodes WHERE city = 'San Francisco' AND zip = '90210';
```

even though there will really be zero rows satisfying this query. Functional dependency statistics do not provide enough information to conclude that, however.

In many practical situations, this assumption is usually satisfied; for example, there might be a GUI in the application that only allows selecting compatible city and ZIP code values to use in a query. But if that's not the case, functional dependencies may not be a viable option.

### 14.2.2.2. Multivariate N-Distinct Counts

Single-column statistics store the number of distinct values in each column. Estimates of the number of distinct values when combining more than one column (for example, for GROUP BY a, b) are frequently wrong when the planner only has single-column statistical data, causing it to select bad plans.

To improve such estimates, ANALYZE can collect n-distinct statistics for groups of columns. As before, it's impractical to do this for every possible column grouping, so data is collected only for those
groups of columns appearing together in a statistics object defined with the ndistinct option. Data will be collected for each possible combination of two or more columns from the set of listed columns.

Continuing the previous example, the n-distinct counts in a table of ZIP codes might look like the following:

```sql
CREATE STATISTICS stts2 (ndistinct) ON zip, state, city FROM zipcodes;
ANALYZE zipcodes;
SELECT stxkeys AS k, stxndistinct AS nd
FROM pg_statistic_ext
WHERE stxname = 'stts2';
```

<table>
<thead>
<tr>
<th>k</th>
<th>nd</th>
</tr>
</thead>
<tbody>
<tr>
<td>1, 2</td>
<td>33178, 33178</td>
</tr>
<tr>
<td>1, 5</td>
<td>33178, 33178</td>
</tr>
<tr>
<td>2, 5</td>
<td>27435, 27435</td>
</tr>
</tbody>
</table>

This indicates that there are three combinations of columns that have 33178 distinct values: ZIP code and state; ZIP code and city; and ZIP code, city and state (the fact that they are all equal is expected given that ZIP code alone is unique in this table). On the other hand, the combination of city and state has only 27435 distinct values.

It’s advisable to create ndistinct statistics objects only on combinations of columns that are actually used for grouping, and for which misestimation of the number of groups is resulting in bad plans. Otherwise, the ANALYZE cycles are just wasted.

### 14.3. Controlling the Planner with Explicit JOIN Clauses

It is possible to control the query planner to some extent by using the explicit JOIN syntax. To see why this matters, we first need some background.

In a simple join query, such as:

```sql
SELECT * FROM a, b, c WHERE a.id = b.id AND b.ref = c.id;
```

the planner is free to join the given tables in any order. For example, it could generate a query plan that joins A to B, using the WHERE condition `a.id = b.id`, and then joins C to this joined table, using the other WHERE condition. Or it could join B to C and then join A to that result. Or it could join A to C and then join them with B — but that would be inefficient, since the full Cartesian product of A and C would have to be formed, there being no applicable condition in the WHERE clause to allow optimization of the join. (All joins in the PostgreSQL executor happen between two input tables, so it’s necessary to build up the result in one or another of these fashions.) The important point is that these different join possibilities give semantically equivalent results but might have hugely different execution costs. Therefore, the planner will explore all of them to try to find the most efficient query plan.

When a query only involves two or three tables, there aren’t many join orders to worry about. But the number of possible join orders grows exponentially as the number of tables expands. Beyond ten or so input tables it’s no longer practical to do an exhaustive search of all the possibilities, and even for six or seven tables planning might take an annoyingly long time. When there are too many input tables, the PostgreSQL planner will switch from exhaustive search to a genetic probabilistic search through
Performance Tips

a limited number of possibilities. (The switch-over threshold is set by the geqo_threshold run-time parameter.) The genetic search takes less time, but it won't necessarily find the best possible plan.

When the query involves outer joins, the planner has less freedom than it does for plain (inner) joins. For example, consider:

```
SELECT * FROM a LEFT JOIN (b JOIN c ON (b.ref = c.id)) ON (a.id = b.id);
```

Although this query's restrictions are superficially similar to the previous example, the semantics are different because a row must be emitted for each row of A that has no matching row in the join of B and C. Therefore the planner has no choice of join order here: it must join B to C and then join A to that result. Accordingly, this query takes less time to plan than the previous query. In other cases, the planner might be able to determine that more than one join order is safe. For example, given:

```
SELECT * FROM a LEFT JOIN b ON (a.bid = b.id) LEFT JOIN c ON (a.cid = c.id);
```

it is valid to join A to either B or C first. Currently, only FULL JOIN completely constrains the join order. Most practical cases involving LEFT JOIN or RIGHT JOIN can be rearranged to some extent.

Explicit inner join syntax (INNER JOIN, CROSS JOIN, or unadorned JOIN) is semantically the same as listing the input relations in FROM, so it does not constrain the join order.

Even though most kinds of JOIN don't completely constrain the join order, it is possible to instruct the PostgreSQL query planner to treat all JOIN clauses as constraining the join order anyway. For example, these three queries are logically equivalent:

```
SELECT * FROM a, b, c WHERE a.id = b.id AND b.ref = c.id;
SELECT * FROM a CROSS JOIN b CROSS JOIN c WHERE a.id = b.id AND b.ref = c.id;
SELECT * FROM a JOIN (b JOIN c ON (b.ref = c.id)) ON (a.id = b.id);
```

But if we tell the planner to honor the JOIN order, the second and third take less time to plan than the first. This effect is not worth worrying about for only three tables, but it can be a lifesaver with many tables.

To force the planner to follow the join order laid out by explicit JOINs, set the join_collapse_limit run-time parameter to 1. (Other possible values are discussed below.)

You do not need to constrain the join order completely in order to cut search time, because it's OK to use JOIN operators within items of a plain FROM list. For example, consider:

```
SELECT * FROM a CROSS JOIN b, c, d, e WHERE ...;
```

With join_collapse_limit = 1, this forces the planner to join A to B before joining them to other tables, but doesn't constrain its choices otherwise. In this example, the number of possible join orders is reduced by a factor of 5.

Constraining the planner's search in this way is a useful technique both for reducing planning time and for directing the planner to a good query plan. If the planner chooses a bad join order by default, you can force it to choose a better order via JOIN syntax — assuming that you know of a better order, that is. Experimentation is recommended.

A closely related issue that affects planning time is collapsing of subqueries into their parent query. For example, consider:
SELECT *
FROM x, y,
    (SELECT * FROM a, b, c WHERE something) AS ss
WHERE somethingelse;

This situation might arise from use of a view that contains a join; the view's SELECT rule will be inserted in place of the view reference, yielding a query much like the above. Normally, the planner will try to collapse the subquery into the parent, yielding:

SELECT * FROM x, y, a, b, c WHERE something AND somethingelse;

This usually results in a better plan than planning the subquery separately. (For example, the outer WHERE conditions might be such that joining X to A first eliminates many rows of A, thus avoiding the need to form the full logical output of the subquery.) But at the same time, we have increased the planning time; here, we have a five-way join problem replacing two separate three-way join problems. Because of the exponential growth of the number of possibilities, this makes a big difference. The planner tries to avoid getting stuck in huge join search problems by not collapsing a subquery if more than from_collapse_limit FROM items would result in the parent query. You can trade off planning time against quality of plan by adjusting this run-time parameter up or down.

from_collapse_limit and join_collapse_limit are similarly named because they do almost the same thing: one controls when the planner will “flatten out” subqueries, and the other controls when it will flatten out explicit joins. Typically you would either set join_collapse_limit equal to from_collapse_limit (so that explicit joins and subqueries act similarly) or set join_collapse_limit to 1 (if you want to control join order with explicit joins). But you might set them differently if you are trying to fine-tune the trade-off between planning time and run time.

14.4. Populating a Database

One might need to insert a large amount of data when first populating a database. This section contains some suggestions on how to make this process as efficient as possible.

14.4.1. Disable Autocommit

When using multiple INSERTs, turn off autocommit and just do one commit at the end. (In plain SQL, this means issuing BEGIN at the start and COMMIT at the end. Some client libraries might do this behind your back, in which case you need to make sure the library does it when you want it done.) If you allow each insertion to be committed separately, PostgreSQL is doing a lot of work for each row that is added. An additional benefit of doing all insertions in one transaction is that if the insertion of one row were to fail then the insertion of all rows inserted up to that point would be rolled back, so you won’t be stuck with partially loaded data.

14.4.2. Use COPY

Use COPY to load all the rows in one command, instead of using a series of INSERT commands. The COPY command is optimized for loading large numbers of rows; it is less flexible than INSERT, but incurs significantly less overhead for large data loads. Since COPY is a single command, there is no need to disable autocommit if you use this method to populate a table.

If you cannot use COPY, it might help to use PREPARE to create a prepared INSERT statement, and then use EXECUTE as many times as required. This avoids some of the overhead of repeatedly parsing and planning INSERT. Different interfaces provide this facility in different ways; look for “prepared statements” in the interface documentation.

Note that loading a large number of rows using COPY is almost always faster than using INSERT, even if PREPARE is used and multiple insertions are batched into a single transaction.
COPY is fastest when used within the same transaction as an earlier CREATE TABLE or TRUNCATE command. In such cases no WAL needs to be written, because in case of an error, the files containing the newly loaded data will be removed anyway. However, this consideration only applies when wal_level is minimal as all commands must write WAL otherwise.

14.4.3. Remove Indexes

If you are loading a freshly created table, the fastest method is to create the table, bulk load the table's data using COPY, then create any indexes needed for the table. Creating an index on pre-existing data is quicker than updating it incrementally as each row is loaded.

If you are adding large amounts of data to an existing table, it might be a win to drop the indexes, load the table, and then recreate the indexes. Of course, the database performance for other users might suffer during the time the indexes are missing. One should also think twice before dropping a unique index, since the error checking afforded by the unique constraint will be lost while the index is missing.

14.4.4. Remove Foreign Key Constraints

Just as with indexes, a foreign key constraint can be checked “in bulk” more efficiently than row-by-row. So it might be useful to drop foreign key constraints, load data, and re-create the constraints. Again, there is a trade-off between data load speed and loss of error checking while the constraint is missing.

What’s more, when you load data into a table with existing foreign key constraints, each new row requires an entry in the server’s list of pending trigger events (since it is the firing of a trigger that checks the row’s foreign key constraint). Loading many millions of rows can cause the trigger event queue to overflow available memory, leading to intolerable swapping or even outright failure of the command. Therefore it may be necessary, not just desirable, to drop and re-apply foreign keys when loading large amounts of data. If temporarily removing the constraint isn’t acceptable, the only other recourse may be to split up the load operation into smaller transactions.

14.4.5. Increase maintenance_work_mem

Temporarily increasing the maintenance_work_mem configuration variable when loading large amounts of data can lead to improved performance. This will help to speed up CREATE INDEX commands and ALTER TABLE ADD FOREIGN KEY commands. It won’t do much for COPY itself, so this advice is only useful when you are using one or both of the above techniques.

14.4.6. Increase max_wal_size

Temporarily increasing the max_wal_size configuration variable can also make large data loads faster. This is because loading a large amount of data into PostgreSQL will cause checkpoints to occur more often than the normal checkpoint frequency (specified by the checkpoint_timeout configuration variable). Whenever a checkpoint occurs, all dirty pages must be flushed to disk. By increasing max_wal_size temporarily during bulk data loads, the number of checkpoints that are required can be reduced.

14.4.7. Disable WAL Archival and Streaming Replication

When loading large amounts of data into an installation that uses WAL archiving or streaming replication, it might be faster to take a new base backup after the load has completed than to process a large amount of incremental WAL data. To prevent incremental WAL logging while loading, disable archiving and streaming replication, by setting wal_level to minimal, archive_mode to off, and max_wal_senders to zero. But note that changing these settings requires a server restart.
Aside from avoiding the time for the archiver or WAL sender to process the WAL data, doing this will actually make certain commands faster, because they are designed not to write WAL at all if `wal_level` is `minimal`. (They can guarantee crash safety more cheaply by doing an `fsync` at the end than by writing WAL.) This applies to the following commands:

- `CREATE TABLE AS SELECT`
- `CREATE INDEX` (and variants such as `ALTER TABLE ADD PRIMARY KEY`)
- `ALTER TABLE SET TABLESPACE`
- `CLUSTER`
- `COPY FROM`, when the target table has been created or truncated earlier in the same transaction

### 14.4.8. Run `ANALYZE` Afterwards

Whenever you have significantly altered the distribution of data within a table, running `ANALYZE` is strongly recommended. This includes bulk loading large amounts of data into the table. Running `ANALYZE` (or `VACUUM ANALYZE`) ensures that the planner has up-to-date statistics about the table. With no statistics or obsolete statistics, the planner might make poor decisions during query planning, leading to poor performance on any tables with inaccurate or nonexistent statistics. Note that if the autovacuum daemon is enabled, it might run `ANALYZE` automatically; see Section 24.1.3 and Section 24.1.6 for more information.

### 14.4.9. Some Notes About `pg_dump`

Dump scripts generated by `pg_dump` automatically apply several, but not all, of the above guidelines. To reload a `pg_dump` dump as quickly as possible, you need to do a few extra things manually. (Note that these points apply while restoring a dump, not while creating it. The same points apply whether loading a text dump with `psql` or using `pg_restore` to load from a `pg_dump` archive file.)

By default, `pg_dump` uses `COPY`, and when it is generating a complete schema-and-data dump, it is careful to load data before creating indexes and foreign keys. So in this case several guidelines are handled automatically. What is left for you to do is to:

- Set appropriate (i.e., larger than normal) values for `maintenance_work_mem` and `max_wal_size`.
- If using WAL archiving or streaming replication, consider disabling them during the restore. To do that, set `archive_mode` to `off`, `wal_level` to `minimal`, and `max_wal_senders` to zero before loading the dump. Afterwards, set them back to the right values and take a fresh base backup.
- Experiment with the parallel dump and restore modes of both `pg_dump` and `pg_restore` and find the optimal number of concurrent jobs to use. Dumping and restoring in parallel by means of the `-j` option should give you a significantly higher performance over the serial mode.
- Consider whether the whole dump should be restored as a single transaction. To do that, pass the `-l` or `--single-transaction` command-line option to `psql` or `pg_restore`. When using this mode, even the smallest of errors will rollback the entire restore, possibly discarding many hours of processing. Depending on how interrelated the data is, that might seem preferable to manual cleanup, or not. `COPY` commands will run fastest if you use a single transaction and have WAL archiving turned off.
- If multiple CPUs are available in the database server, consider using `pg_restore`'s `--jobs` option. This allows concurrent data loading and index creation.
- Run `ANALYZE` afterwards.
A data-only dump will still use COPY, but it does not drop or recreate indexes, and it does not normally touch foreign keys. So when loading a data-only dump, it is up to you to drop and recreate indexes and foreign keys if you wish to use those techniques. It’s still useful to increase max_wal_size while loading the data, but don’t bother increasing maintenance_work_mem; rather, you’d do that while manually recreating indexes and foreign keys afterwards. And don’t forget to ANALYZE when you’re done; see Section 24.1.3 and Section 24.1.6 for more information.

14.5. Non-Durable Settings

Durability is a database feature that guarantees the recording of committed transactions even if the server crashes or loses power. However, durability adds significant database overhead, so if your site does not require such a guarantee, PostgreSQL can be configured to run much faster. The following are configuration changes you can make to improve performance in such cases. Except as noted below, durability is still guaranteed in case of a crash of the database software; only abrupt operating system stoppage creates a risk of data loss or corruption when these settings are used.

• Place the database cluster’s data directory in a memory-backed file system (i.e. RAM disk). This eliminates all database disk I/O, but limits data storage to the amount of available memory (and perhaps swap).

• Turn off fsync; there is no need to flush data to disk.

• Turn off synchronous_commit; there might be no need to force WAL writes to disk on every commit. This setting does risk transaction loss (though not data corruption) in case of a crash of the database.

• Turn off full_page_writes; there is no need to guard against partial page writes.

• Increase max_wal_size and checkpoint_timeout; this reduces the frequency of checkpoints, but increases the storage requirements of /pg_wal.

• Create unlogged tables to avoid WAL writes, though it makes the tables non-crash-safe.

---

1 You can get the effect of disabling foreign keys by using the --disable-triggers option — but realize that that eliminates, rather than just postpones, foreign key validation, and so it is possible to insert bad data if you use it.
Chapter 15. Parallel Query

PostgreSQL can devise query plans which can leverage multiple CPUs in order to answer queries faster. This feature is known as parallel query. Many queries cannot benefit from parallel query, either due to limitations of the current implementation or because there is no imaginable query plan which is any faster than the serial query plan. However, for queries that can benefit, the speedup from parallel query is often very significant. Many queries can run more than twice as fast when using parallel query, and some queries can run four times faster or even more. Queries that touch a large amount of data but return only a few rows to the user will typically benefit most. This chapter explains some details of how parallel query works and in which situations it can be used so that users who wish to make use of it can understand what to expect.

15.1. How Parallel Query Works

When the optimizer determines that parallel query is the fastest execution strategy for a particular query, it will create a query plan which includes a Gather or Gather Merge node. Here is a simple example:

```sql
EXPLAIN SELECT * FROM pgbench_accounts WHERE filler LIKE '%x%';
```

```
QUERY PLAN
-------------------------------------------------------------------------------------
Gather  (cost=1000.00..217018.43 rows=1 width=97)
  Workers Planned: 2
   ->  Parallel Seq Scan on pgbench_accounts  (cost=0.00..216018.33 rows=1 width=97)
       Filter: (filler ~~ '%x%'::text)
(4 rows)
```

In all cases, the Gather or Gather Merge node will have exactly one child plan, which is the portion of the plan that will be executed in parallel. If the Gather or Gather Merge node is at the very top of the plan tree, then the entire query will execute in parallel. If it is somewhere else in the plan tree, then only the portion of the plan below it will run in parallel. In the example above, the query accesses only one table, so there is only one plan node other than the Gather node itself; since that plan node is a child of the Gather node, it will run in parallel.

Using EXPLAIN, you can see the number of workers chosen by the planner. When the Gather node is reached during query execution, the process which is implementing the user's session will request a number of background worker processes equal to the number of workers chosen by the planner. The number of background workers that the planner will consider using is limited to at most max_parallel_workers_per_gather. The total number of background workers that can exist at any one time is limited by both max_worker_processes and max_parallel_workers. Therefore, it is possible for a parallel query to run with fewer workers than planned, or even with no workers at all. The optimal plan may depend on the number of workers that are available, so this can result in poor query performance. If this occurrence is frequent, consider increasing max_worker_processes and max_parallel_workers so that more workers can be run simultaneously or alternatively reducing max_parallel_workers_per_gather so that the planner requests fewer workers.

Every background worker process which is successfully started for a given parallel query will execute the parallel portion of the plan. The leader will also execute that portion of the plan, but it has an additional responsibility: it must also read all of the tuples generated by the workers. When the parallel portion of the plan generates only a small number of tuples, the leader will often behave very much like an additional worker, speeding up query execution. Conversely, when the parallel portion of the plan generates a large number of tuples, the leader may be almost entirely occupied with reading the tuples generated by the workers and performing any further processing steps which are required by
plan nodes above the level of the \texttt{Gather} node or \texttt{Gather Merge} node. In such cases, the leader will do very little of the work of executing the parallel portion of the plan.

When the node at the top of the parallel portion of the plan is \texttt{Gather Merge} rather than \texttt{Gather}, it indicates that each process executing the parallel portion of the plan is producing tuples in sorted order, and that the leader is performing an order-preserving merge. In contrast, \texttt{Gather} reads tuples from the workers in whatever order is convenient, destroying any sort order that may have existed.

## 15.2. When Can Parallel Query Be Used?

There are several settings which can cause the query planner not to generate a parallel query plan under any circumstances. In order for any parallel query plans whatsoever to be generated, the following settings must be configured as indicated.

- \texttt{max_parallel_workers_per_gather} must be set to a value which is greater than zero. This is a special case of the more general principle that no more workers should be used than the number configured via \texttt{max_parallel_workers_per_gather}.

- \texttt{dynamic_shared_memory_type} must be set to a value other than \texttt{none}. Parallel query requires dynamic shared memory in order to pass data between cooperating processes.

In addition, the system must not be running in single-user mode. Since the entire database system is running in single process in this situation, no background workers will be available.

Even when it is in general possible for parallel query plans to be generated, the planner will not generate them for a given query if any of the following are true:

- The query writes any data or locks any database rows. If a query contains a data-modifying operation either at the top level or within a CTE, no parallel plans for that query will be generated. As an exception, the commands \texttt{CREATE TABLE ... AS}, \texttt{SELECT INTO}, and \texttt{CREATE MATERIALIZED VIEW} which create a new table and populate it can use a parallel plan.

- The query might be suspended during execution. In any situation in which the system thinks that partial or incremental execution might occur, no parallel plan is generated. For example, a cursor created using \texttt{DECLARE CURSOR} will never use a parallel plan. Similarly, a PL/pgSQL loop of the form \texttt{FOR x IN query LOOP .. END LOOP} will never use a parallel plan, because the parallel query system is unable to verify that the code in the loop is safe to execute while parallel query is active.

- The query uses any function marked \texttt{PARALLEL UNSAFE}. Most system-defined functions are \texttt{PARALLEL SAFE}, but user-defined functions are marked \texttt{PARALLEL UNSAFE} by default. See the discussion of Section 15.4.

- The query is running inside of another query that is already parallel. For example, if a function called by a parallel query issues an SQL query itself, that query will never use a parallel plan. This is a limitation of the current implementation, but it may not be desirable to remove this limitation, since it could result in a single query using a very large number of processes.

- The transaction isolation level is serializable. This is a limitation of the current implementation.

Even when parallel query plan is generated for a particular query, there are several circumstances under which it will be impossible to execute that plan in parallel at execution time. If this occurs, the leader will execute the portion of the plan below the \texttt{Gather} node entirely by itself, almost as if the \texttt{Gather} node were not present. This will happen if any of the following conditions are met:

- No background workers can be obtained because of the limitation that the total number of background workers cannot exceed \texttt{max_worker_processes}.

- No background workers can be obtained because of the limitation that the total number of background workers launched for purposes of parallel query cannot exceed \texttt{max_parallel_workers}.
• The client sends an Execute message with a non-zero fetch count. See the discussion of the extended query protocol. Since libpq currently provides no way to send such a message, this can only occur when using a client that does not rely on libpq. If this is a frequent occurrence, it may be a good idea to set `max_parallel_workers_per_gather` to zero in sessions where it is likely, so as to avoid generating query plans that may be suboptimal when run serially.

• The transaction isolation level is serializable. This situation does not normally arise, because parallel query plans are not generated when the transaction isolation level is serializable. However, it can happen if the transaction isolation level is changed to serializable after the plan is generated and before it is executed.

15.3. Parallel Plans

Because each worker executes the parallel portion of the plan to completion, it is not possible to simply take an ordinary query plan and run it using multiple workers. Each worker would produce a full copy of the output result set, so the query would not run any faster than normal but would produce incorrect results. Instead, the parallel portion of the plan must be what is known internally to the query optimizer as a `partial plan`; that is, it must be constructed so that each process which executes the plan will generate only a subset of the output rows in such a way that each required output row is guaranteed to be generated by exactly one of the cooperating processes. Generally, this means that the scan on the driving table of the query must be a parallel-aware scan.

15.3.1. Parallel Scans

The following types of parallel-aware table scans are currently supported.

• In a `parallel sequential scan`, the table’s blocks will be divided among the cooperating processes. Blocks are handed out one at a time, so that access to the table remains sequential.

• In a `parallel bitmap heap scan`, one process is chosen as the leader. That process performs a scan of one or more indexes and builds a bitmap indicating which table blocks need to be visited. These blocks are then divided among the cooperating processes as in a parallel sequential scan. In other words, the heap scan is performed in parallel, but the underlying index scan is not.

• In a `parallel index scan` or `parallel index-only scan`, the cooperating processes take turns reading data from the index. Currently, parallel index scans are supported only for btree indexes. Each process will claim a single index block and will scan and return all tuples referenced by that block; other process can at the same time be returning tuples from a different index block. The results of a parallel btree scan are returned in sorted order within each worker process.

Other scan types, such as scans of non-btree indexes, may support parallel scans in the future.

15.3.2. Parallel Joins

Just as in a non-parallel plan, the driving table may be joined to one or more other tables using a nested loop, hash join, or merge join. The inner side of the join may be any kind of non-parallel plan that is otherwise supported by the planner provided that it is safe to run within a parallel worker. Depending on the join type, the inner side may also be a parallel plan.

• In a `nested loop join`, the inner side is always non-parallel. Although it is executed in full, this is efficient if the inner side is an index scan, because the outer tuples and thus the loops that look up values in the index are divided over the cooperating processes.

• In a `merge join`, the inner side is always a non-parallel plan and therefore executed in full. This may be inefficient, especially if a sort must be performed, because the work and resulting data are duplicated in every cooperating process.

• In a `hash join` (without the "parallel" prefix), the inner side is executed in full by every cooperating process to build identical copies of the hash table. This may be inefficient if the hash table is large
or the plan is expensive. In a parallel hash join, the inner side is a parallel hash that divides the work of building a shared hash table over the cooperating processes.

15.3.3. Parallel Aggregation

PostgreSQL supports parallel aggregation by aggregating in two stages. First, each process participating in the parallel portion of the query performs an aggregation step, producing a partial result for each group of which that process is aware. This is reflected in the plan as a Partial Aggregate node. Second, the partial results are transferred to the leader via Gather or Gather Merge. Finally, the leader re-aggregates the results across all workers in order to produce the final result. This is reflected in the plan as a Finalize Aggregate node.

Because the Finalize Aggregate node runs on the leader process, queries which produce a relatively large number of groups in comparison to the number of input rows will appear less favorable to the query planner. For example, in the worst-case scenario the number of groups seen by the Finalize Aggregate node could be as many as the number of input rows which were seen by all worker processes in the Partial Aggregate stage. For such cases, there is clearly going to be no performance benefit to using parallel aggregation. The query planner takes this into account during the planning process and is unlikely to choose parallel aggregate in this scenario.

Parallel aggregation is not supported in all situations. Each aggregate must be safe for parallelism and must have a combine function. If the aggregate has a transition state of type internal, it must have serialization and deserialization functions. See CREATE AGGREGATE for more details. Parallel aggregation is not supported if any aggregate function call contains DISTINCT or ORDER BY clause and is also not supported for ordered set aggregates or when the query involves GROUPING SETS. It can only be used when all joins involved in the query are also part of the parallel portion of the plan.

15.3.4. Parallel Append

Whenever PostgreSQL needs to combine rows from multiple sources into a single result set, it uses an Append or MergeAppend plan node. This commonly happens when implementing UNION ALL or when scanning a partitioned table. Such nodes can be used in parallel plans just as they can in any other plan. However, in a parallel plan, the planner may instead use a Parallel Append node.

When an Append node is used in a parallel plan, each process will execute the child plans in the order in which they appear, so that all participating processes cooperate to execute the first child plan until it is complete and then move to the second plan at around the same time. When a Parallel Append is used instead, the executor will instead spread out the participating processes as evenly as possible across its child plans, so that multiple child plans are executed simultaneously. This avoids contention, and also avoids paying the startup cost of a child plan in those processes that never execute it.

Also, unlike a regular Append node, which can only have partial children when used within a parallel plan, a Parallel Append node can have both partial and non-partial child plans. Non-partial children will be scanned by only a single process, since scanning them more than once would produce duplicate results. Plans that involve appending multiple results sets can therefore achieve coarse-grained parallelism even when efficient partial plans are not available. For example, consider a query against a partitioned table which can be only be implemented efficiently by using an index that does not support parallel scans. The planner might choose a Parallel Append of regular Index Scan plans; each individual index scan would have to be executed to completion by a single process, but different scans could be performed at the same time by different processes.

enable_parallel_append can be used to disable this feature.

15.3.5. Parallel Plan Tips

If a query that is expected to do so does not produce a parallel plan, you can try reducing parallel_set-up_cost or parallel_tuple_cost. Of course, this plan may turn out to be slower than the serial plan which the planner preferred, but this will not always be the case. If you don't get a parallel plan even with very
small values of these settings (e.g. after setting them both to zero), there may be some reason why the query planner is unable to generate a parallel plan for your query. See Section 15.2 and Section 15.4 for information on why this may be the case.

When executing a parallel plan, you can use `EXPLAIN (ANALYZE, VERBOSE)` to display per-worker statistics for each plan node. This may be useful in determining whether the work is being evenly distributed between all plan nodes and more generally in understanding the performance characteristics of the plan.

### 15.4. Parallel Safety

The planner classifies operations involved in a query as either parallel safe, parallel restricted, or parallel unsafe. A parallel safe operation is one which does not conflict with the use of parallel query. A parallel restricted operation is one which cannot be performed in a parallel worker, but which can be performed in the leader while parallel query is in use. Therefore, parallel restricted operations can never occur below a Gather or Gather Merge node, but can occur elsewhere in a plan which contains such a node. A parallel unsafe operation is one which cannot be performed while parallel query is in use, not even in the leader. When a query contains anything which is parallel unsafe, parallel query is completely disabled for that query.

The following operations are always parallel restricted.

- Scans of common table expressions (CTEs).
- Scans of temporary tables.
- Scans of foreign tables, unless the foreign data wrapper has an `IsForeignScanParallelSafe` API which indicates otherwise.
- Plan nodes to which an `InitPlan` is attached.
- Plan nodes which reference a correlated `SubPlan`.

#### 15.4.1. Parallel Labeling for Functions and Aggregates

The planner cannot automatically determine whether a user-defined function or aggregate is parallel safe, parallel restricted, or parallel unsafe, because this would require predicting every operation which the function could possibly perform. In general, this is equivalent to the Halting Problem and therefore impossible. Even for simple functions where it could conceivably be done, we do not try, since this would be expensive and error-prone. Instead, all user-defined functions are assumed to be parallel unsafe unless otherwise marked. When using `CREATE FUNCTION` or `ALTER FUNCTION`, markings can be set by specifying `PARALLEL SAFE`, `PARALLEL RESTRICTED`, or `PARALLEL UNSAFE` as appropriate. When using `CREATE AGGREGATE`, the `PARALLEL` option can be specified with `SAFE`, `RESTRICTED`, or `UNSAFE` as the corresponding value.

Functions and aggregates must be marked `PARALLEL UNSAFE` if they write to the database, access sequences, change the transaction state even temporarily (e.g. a PL/pgSQL function which establishes an `EXCEPTION` block to catch errors), or make persistent changes to settings. Similarly, functions must be marked `PARALLEL RESTRICTED` if they access temporary tables, client connection state, cursors, prepared statements, or miscellaneous backend-local state which the system cannot synchronize across workers. For example, `setseed` and `random` are parallel restricted for this last reason.

In general, if a function is labeled as being safe when it is restricted or unsafe, or if it is labeled as being restricted when it is in fact unsafe, it may throw errors or produce wrong answers when used in a parallel query. C-language functions could in theory exhibit totally undefined behavior if mislabeled, since there is no way for the system to protect itself against arbitrary C code, but in most likely cases the result will be no worse than for any other function. If in doubt, it is probably best to label functions as `UNSAFE`.
Parallel Query

If a function executed within a parallel worker acquires locks which are not held by the leader, for example by querying a table not referenced in the query, those locks will be released at worker exit, not end of transaction. If you write a function which does this, and this behavior difference is important to you, mark such functions as **PARALLEL RESTRICTED** to ensure that they execute only in the leader.

Note that the query planner does not consider deferring the evaluation of parallel-restricted functions or aggregates involved in the query in order to obtain a superior plan. So, for example, if a **WHERE** clause applied to a particular table is parallel restricted, the query planner will not consider performing a scan of that table in the parallel portion of a plan. In some cases, it would be possible (and perhaps even efficient) to include the scan of that table in the parallel portion of the query and defer the evaluation of the **WHERE** clause so that it happens above the **Gather** node. However, the planner does not do this.
Part III. Server Administration

This part covers topics that are of interest to a PostgreSQL database administrator. This includes installation of the software, set up and configuration of the server, management of users and databases, and maintenance tasks. Anyone who runs a PostgreSQL server, even for personal use, but especially in production, should be familiar with the topics covered in this part.

The information in this part is arranged approximately in the order in which a new user should read it. But the chapters are self-contained and can be read individually as desired. The information in this part is presented in a narrative fashion in topical units. Readers looking for a complete description of a particular command should see Part VI.

The first few chapters are written so they can be understood without prerequisite knowledge, so new users who need to set up their own server can begin their exploration with this part. The rest of this part is about tuning and management; that material assumes that the reader is familiar with the general use of the PostgreSQL database system. Readers are encouraged to look at Part I and Part II for additional information.
# Table of Contents

16. Installation from Source Code .............................................................. 474
   16.1. Short Version ............................................................................. 474
   16.2. Requirements ........................................................................... 474
   16.3. Getting The Source .................................................................... 476
   16.4. Installation Procedure ................................................................. 476
   16.5. Post-Installation Setup ................................................................. 488
       16.5.1. Shared Libraries ................................................................. 488
       16.5.2. Environment Variables ....................................................... 489
   16.6. Supported Platforms ................................................................. 490
   16.7. Platform-specific Notes ............................................................... 490
       16.7.1. AIX ................................................................................ 490
       16.7.2. Cygwin .......................................................................... 493
       16.7.3. HP-UX .......................................................................... 494
       16.7.4. macOS ......................................................................... 495
       16.7.5. MinGW/Native Windows ..................................................... 495
       16.7.6. Solaris ............................................................................ 496

17. Installation from Source Code on Windows ........................................... 498
   17.1. Building with Visual C++ or the Microsoft Windows SDK .......... 498
       17.1.1. Requirements ................................................................... 499
       17.1.2. Special Considerations for 64-bit Windows ......................... 501
       17.1.3. Building .......................................................................... 501
       17.1.4. Cleaning and Installing ...................................................... 501
       17.1.5. Running the Regression Tests ............................................. 502
       17.1.6. Building the Documentation ............................................. 502

18. Server Setup and Operation ................................................................. 504
   18.1. The PostgreSQL User Account .................................................... 504
   18.2. Creating a Database Cluster ......................................................... 504
       18.2.1. Use of Secondary File Systems ........................................... 505
       18.2.2. Use of Network File Systems ............................................. 506
   18.3. Starting the Database Server ....................................................... 506
       18.3.1. Server Start-up Failures ...................................................... 508
       18.3.2. Client Connection Problems ............................................. 509
   18.4. Managing Kernel Resources ....................................................... 509
       18.4.1. Shared Memory and Semaphores ..................................... 509
       18.4.2. systemd RemoveIPC ........................................................ 514
       18.4.3. Resource Limits ............................................................... 515
       18.4.4. Linux Memory Overcommit .............................................. 516
       18.4.5. Linux Huge Pages ............................................................ 517
   18.5. Shutting Down the Server ........................................................... 518
   18.6. Upgrading a PostgreSQL Cluster ............................................... 519
       18.6.1. Upgrading Data via pg_dumpall ....................................... 520
       18.6.2. Upgrading Data via pg_upgrade ....................................... 521
       18.6.3. Upgrading Data via Replication ....................................... 521
   18.7. Preventing Server Spoofing ........................................................ 522
   18.8. Encryption Options ................................................................. 522
   18.9. Secure TCP/IP Connections with SSL ........................................ 523
       18.9.1. Basic Setup .................................................................... 523
       18.9.2. OpenSSL Configuration .................................................... 524
       18.9.3. Using Client Certificates .................................................. 524
       18.9.4. SSL Server File Usage ..................................................... 525
       18.9.5. Creating Certificates ....................................................... 525
   18.10. Secure TCP/IP Connections with SSH Tunnels ......................... 527
   18.11. Registering Event Log on Windows ........................................... 528

19. Server Configuration ........................................................................ 529
   19.1. Setting Parameters ..................................................................... 529
19.1. Parameter Names and Values ................................................................. 529
19.1.2. Parameter Interaction via the Configuration File .......................... 529
19.1.3. Parameter Interaction via SQL ......................................................... 530
19.1.4. Parameter Interaction via the Shell ............................................... 531
19.1.5. Managing Configuration File Contents ......................................... 531
19.2. File Locations ...................................................................................... 532
19.3. Connections and Authentication ......................................................... 533
  19.3.1. Connection Settings ...................................................................... 533
  19.3.2. Authentication .............................................................................. 536
  19.3.3. SSL ............................................................................................... 537
19.4. Resource Consumption ........................................................................ 539
  19.4.1. Memory ......................................................................................... 539
  19.4.2. Disk ............................................................................................... 541
  19.4.3. Kernel Resource Usage ................................................................. 541
  19.4.4. Cost-based Vacuum Delay ........................................................... 541
  19.4.5. Background Writer ....................................................................... 542
  19.4.6. Asynchronous Behavior ............................................................... 543
19.5. Write Ahead Log .................................................................................. 545
  19.5.1. Settings .......................................................................................... 545
  19.5.2. Checkpoints ............................................................................... 549
  19.5.3. Archiving ...................................................................................... 550
19.6. Replication ............................................................................................ 551
  19.6.1. Sending Servers .......................................................................... 551
  19.6.2. Master Server ............................................................................. 552
  19.6.3. Standby Servers .......................................................................... 554
  19.6.4. Subscribers .................................................................................. 555
19.7. Query Planning .................................................................................... 556
  19.7.1. Planner Method Configuration ...................................................... 556
  19.7.2. Planner Cost Constants ................................................................. 557
  19.7.3. Genetic Query Optimizer ............................................................... 559
  19.7.4. Other Planner Options ................................................................. 560
19.8. Error Reporting and Logging ............................................................. 562
  19.8.1. Where To Log .............................................................................. 562
  19.8.2. When To Log ............................................................................... 566
  19.8.3. What To Log .............................................................................. 567
  19.8.4. Using CSV-Format Log Output .................................................... 571
  19.8.5. Process Title ............................................................................... 572
19.9. Run-time Statistics ............................................................................... 572
  19.9.1. Query and Index Statistics Collector ............................................ 572
  19.9.2. Statistics Monitoring ................................................................. 573
19.10. Automatic Vacuuming ....................................................................... 574
19.11. Client Connection Defaults .............................................................. 575
  19.11.1. Statement Behavior ................................................................... 575
  19.11.2. Locale and Formatting ............................................................... 580
  19.11.3. Shared Library Preloading ......................................................... 582
  19.11.4. Other Defaults ........................................................................... 583
19.12. Lock Management ............................................................................. 584
19.13. Version and Platform Compatibility ................................................ 585
  19.13.1. Previous PostgreSQL Versions ................................................... 585
  19.13.2. Platform and Client Compatibility ............................................ 586
19.15. Preset Options ................................................................................... 587
19.16. Customized Options ......................................................................... 589
19.17. Developer Options ........................................................................... 589
19.18. Short Options .................................................................................... 592
20. Client Authentication .............................................................................. 594
  20.1. The pg_hba.conf File ....................................................................... 594
  20.2. User Name Maps ............................................................................. 601
Chapter 16. Installation from Source Code

This chapter describes the installation of PostgreSQL using the source code distribution. (If you are installing a pre-packaged distribution, such as an RPM or Debian package, ignore this chapter and read the packager's instructions instead.)

16.1. Short Version

```plaintext
cmake
make
su
make install
adduser postgres
mkdir /usr/local/pgsql/data
chown postgres /usr/local/pgsql/data
su - postgres
/usr/local/pgsql/bin/initdb -D /usr/local/pgsql/data
/usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data >logfile
2>&1 &
/usr/local/pgsql/bin/createdb test
/usr/local/pgsql/bin/psql test
```

The long version is the rest of this chapter.

16.2. Requirements

In general, a modern Unix-compatible platform should be able to run PostgreSQL. The platforms that had received specific testing at the time of release are listed in Section 16.6 below. In the doc subdirectory of the distribution there are several platform-specific FAQ documents you might wish to consult if you are having trouble.

The following software packages are required for building PostgreSQL:

- GNU make version 3.80 or newer is required; other make programs or older GNU make versions will not work. (GNU make is sometimes installed under the name gmake.) To test for GNU make enter:

  ```plaintext
  make --version
  ```

- You need an ISO/ANSI C compiler (at least C89-compliant). Recent versions of GCC are recommended, but PostgreSQL is known to build using a wide variety of compilers from different vendors.

- tar is required to unpack the source distribution, in addition to either gzip or bzip2.

- The GNU Readline library is used by default. It allows psql (the PostgreSQL command line SQL interpreter) to remember each command you type, and allows you to use arrow keys to recall and edit previous commands. This is very helpful and is strongly recommended. If you don't want to use it then you must specify the --without-readline option to configure. As an alternative, you can often use the BSD-licensed libedit library, originally developed on NetBSD. The libedit library is GNU Readline-compatible and is used if libreadline is not found, or if --with-libedit-preferred is used as an option to configure. If you are using a package-based Linux distribution, be aware that you need both the readline and readline-devel packages, if those are separate in your distribution.
• The zlib compression library is used by default. If you don't want to use it then you must specify the `--without-zlib` option to `configure`. Using this option disables support for compressed archives in `pg_dump` and `pg_restore`.

The following packages are optional. They are not required in the default configuration, but they are needed when certain build options are enabled, as explained below:

• To build the server programming language PL/Perl you need a full Perl installation, including the `libperl` library and the header files. The minimum required version is Perl 5.8.3. Since PL/Perl will be a shared library, the `libperl` library must be a shared library also on most platforms. This appears to be the default in recent Perl versions, but it was not in earlier versions, and in any case it is the choice of whoever installed Perl at your site. `configure` will fail if building PL/Perl is selected but it cannot find a shared `libperl`. In that case, you will have to rebuild and install Perl manually to be able to build PL/Perl. During the configuration process for Perl, request a shared library.

If you intend to make more than incidental use of PL/Perl, you should ensure that the Perl installation was built with the `usemultiplicity` option enabled (`perl -V` will show whether this is the case).

• To build the PL/Python server programming language, you need a Python installation with the header files and the `distutils` module. The minimum required version is Python 2.4. Python 3 is supported if it's version 3.1 or later; but see Section 46.1 when using Python 3.

Since PL/Python will be a shared library, the `libpython` library must be a shared library also on most platforms. This is not the case in a default Python installation built from source, but a shared library is available in many operating system distributions. `configure` will fail if building PL/Python is selected but it cannot find a shared `libpython`. That might mean that you either have to install additional packages or rebuild (part of) your Python installation to provide this shared library. When building from source, run Python's `configure` with the `--enable-shared` flag.

• To build the PL/Tcl procedural language, you of course need a Tcl installation. The minimum required version is Tcl 8.4.

• To enable Native Language Support (NLS), that is, the ability to display a program's messages in a language other than English, you need an implementation of the Gettext API. Some operating systems have this built-in (e.g., Linux, NetBSD, Solaris), for other systems you can download an add-on package from http://www.gnu.org/software/gettext/. If you are using the Gettext implementation in the GNU C library then you will additionally need the GNU Gettext package for some utility programs. For any of the other implementations you will not need it.

• You need OpenSSL, if you want to support encrypted client connections. The minimum required version is 0.9.8.

• You need Kerberos, OpenLDAP, and/or PAM, if you want to support authentication using those services.

• To build the PostgreSQL documentation, there is a separate set of requirements; see Section J.2.

If you are building from a Git tree instead of using a released source package, or if you want to do server development, you also need the following packages:

• Flex and Bison are needed to build from a Git checkout, or if you changed the actual scanner and parser definition files. If you need them, be sure to get Flex 2.5.31 or later and Bison 1.875 or later. Other lex and yacc programs cannot be used.

• Perl 5.8.3 or later is needed to build from a Git checkout, or if you changed the input files for any of the build steps that use Perl scripts. If building on Windows you will need Perl in any case. Perl is also required to run some test suites.

If you need to get a GNU package, you can find it at your local GNU mirror site (see https://www.gnu.org/prep/ftp for a list) or at ftp://ftp.gnu.org/gnu/.
Also check that you have sufficient disk space. You will need about 100 MB for the source tree during compilation and about 20 MB for the installation directory. An empty database cluster takes about 35 MB; databases take about five times the amount of space that a flat text file with the same data would take. If you are going to run the regression tests you will temporarily need up to an extra 150 MB. Use the df command to check free disk space.

16.3. Getting The Source

The PostgreSQL 11.1 sources can be obtained from the download section of our website: https://www.postgresql.org/download/. You should get a file named postgresql-11.1.tar.gz or postgresql-11.1.tar.bz2. After you have obtained the file, unpack it:

```
gunzip postgresql-11.1.tar.gz
tar xf postgresql-11.1.tar
```

(Use bunzip2 instead of gunzip if you have the .bz2 file.) This will create a directory postgresql-11.1 under the current directory with the PostgreSQL sources. Change into that directory for the rest of the installation procedure.

You can also get the source directly from the version control repository, see Appendix I.

16.4. Installation Procedure

1. **Configuration**

   The first step of the installation procedure is to configure the source tree for your system and choose the options you would like. This is done by running the configure script. For a default installation simply enter:

   ```
   ./configure
   ```

   This script will run a number of tests to determine values for various system dependent variables and detect any quirks of your operating system, and finally will create several files in the build tree to record what it found. You can also run configure in a directory outside the source tree, if you want to keep the build directory separate. This procedure is also called a VPATH build. Here’s how:

   ```
   mkdir build_dir
cd build_dir
/path/to/source/tree/configure [options go here]
make
   ```

   The default configuration will build the server and utilities, as well as all client applications and interfaces that require only a C compiler. All files will be installed under /usr/local/pgsql by default.

   You can customize the build and installation process by supplying one or more of the following command line options to configure:

   ```
   --prefix=PREFIX
   ```

   Install all files under the directory PREFIX instead of /usr/local/pgsql. The actual files will be installed into various subdirectories; no files will ever be installed directly into the PREFIX directory.
If you have special needs, you can also customize the individual subdirectories with the following options. However, if you leave these with their defaults, the installation will be relocatable, meaning you can move the directory after installation. (The man and doc locations are not affected by this.)

For relocatable installs, you might want to use `configure`'s `--disable-rpath` option. Also, you will need to tell the operating system how to find the shared libraries.

`--exec-prefix=EXEC-PREFIX`

You can install architecture-dependent files under a different prefix, `EXEC-PREFIX`, than what `PREFIX` was set to. This can be useful to share architecture-independent files between hosts. If you omit this, then `EXEC-PREFIX` is set equal to `PREFIX` and both architecture-dependent and independent files will be installed under the same tree, which is probably what you want.

`--bindir=DIRECTORY`

Specifies the directory for executable programs. The default is `EXEC-PREFIX/bin`, which normally means `/usr/local/pgsql/bin`.

`--sysconfdir=DIRECTORY`

Sets the directory for various configuration files, `PREFIX/etc` by default.

`--libdir=DIRECTORY`

Sets the location to install libraries and dynamically loadable modules. The default is `EXEC-PREFIX/lib`.

`--includedir=DIRECTORY`

Sets the directory for installing C and C++ header files. The default is `PREFIX/include`.

`--datarootdir=DIRECTORY`

Sets the root directory for various types of read-only data files. This only sets the default for some of the following options. The default is `PREFIX/share`.

`--datadir=DIRECTORY`

Sets the directory for read-only data files used by the installed programs. The default is `DATAROOTDIR`. Note that this has nothing to do with where your database files will be placed.

`--localedir=DIRECTORY`

Sets the directory for installing locale data, in particular message translation catalog files. The default is `DATAROOTDIR/locale`.

`--mandir=DIRECTORY`

The man pages that come with PostgreSQL will be installed under this directory, in their respective `man` subdirectories. The default is `DATAROOTDIR/man`.

`--docdir=DIRECTORY`

Sets the root directory for installing documentation files, except “man” pages. This only sets the default for the following options. The default value for this option is `DATAROOTDIR/doc/postgresql`. 

477
Installation from Source Code

--htmldir=DIRECTORY

The HTML-formatted documentation for PostgreSQL will be installed under this directory. The default is DATAROOTDIR.

Note

Care has been taken to make it possible to install PostgreSQL into shared installation locations (such as /usr/local/include) without interfering with the namespace of the rest of the system. First, the string "/postgresql" is automatically appended to datadir, sysconfdir, and docdir, unless the fully expanded directory name already contains the string "postgres" or "pgsql". For example, if you choose /usr/local as prefix, the documentation will be installed in /usr/local/doc/postgresql, but if the prefix is /opt/postgres, then it will be in /opt/postgres/doc. The public C header files of the client interfaces are installed into includedir and are namespace-clean. The internal header files and the server header files are installed into private directories under includedir. See the documentation of each interface for information about how to access its header files. Finally, a private subdirectory will also be created, if appropriate, under libdir for dynamically loadable modules.

--with-extra-version=STRING

Append STRING to the PostgreSQL version number. You can use this, for example, to mark binaries built from unreleased Git snapshots or containing custom patches with an extra version string such as a git describe identifier or a distribution package release number.

--with-includes=DIRECTORIES

DIRECTORIES is a colon-separated list of directories that will be added to the list the compiler searches for header files. If you have optional packages (such as GNU Readline) installed in a non-standard location, you have to use this option and probably also the corresponding --with-libraries option.

Example: --with-includes=/opt/gnu/include:/usr/sup/include.

--with-libraries=DIRECTORIES

DIRECTORIES is a colon-separated list of directories to search for libraries. You will probably have to use this option (and the corresponding --with-includes option) if you have packages installed in non-standard locations.

Example: --with-libraries=/opt/gnu/lib:/usr/sup/lib.

--enable-nls[=LANGUAGES]

Enables Native Language Support (NLS), that is, the ability to display a program's messages in a language other than English. LANGUAGES is an optional space-separated list of codes of the languages that you want supported, for example --enable-nls='de fr'. (The intersection between your list and the set of actually provided translations will be computed automatically.) If you do not specify a list, then all available translations are installed.

To use this option, you will need an implementation of the Gettext API; see above.
Installation from Source Code

--with-pgport=NUMBER

Set NUMBER as the default port number for server and clients. The default is 5432. The port can always be changed later on, but if you specify it here then both server and clients will have the same default compiled in, which can be very convenient. Usually the only good reason to select a non-default value is if you intend to run multiple PostgreSQL servers on the same machine.

--with-perl

Build the PL/Perl server-side language.

--with-python

Build the PL/Python server-side language.

--with-tcl

Build the PL/Tcl server-side language.

--with-tclconfig=DIRECTORY

Tcl installs the file tclConfig.sh, which contains configuration information needed to build modules interfacing to Tcl. This file is normally found automatically at a well-known location, but if you want to use a different version of Tcl you can specify the directory in which to look for it.

--with-gssapi

Build with support for GSSAPI authentication. On many systems, the GSSAPI (usually a part of the Kerberos installation) system is not installed in a location that is searched by default (e.g., /usr/include, /usr/lib), so you must use the options --with-includes and --with-libraries in addition to this option. configure will check for the required header files and libraries to make sure that your GSSAPI installation is sufficient before proceeding.

--with-krb-srvnam=NAME

The default name of the Kerberos service principal used by GSSAPI. postgres is the default. There's usually no reason to change this unless you have a Windows environment, in which case it must be set to upper case POSTGRES.

--with-llvm

Build with support for LLVM based JIT compilation (see Chapter 32). This requires the LLVM library to be installed. The minimum required version of LLVM is currently 3.9.

llvm-config will be used to find the required compilation options. llvm-config, and then llvm-config-$major-$minor for all supported versions, will be searched on PATH. If that would not yield the correct binary, use LLVM_CONFIG to specify a path to the correct llvm-config. For example

```
./configure ... --with-llvm LLVM_CONFIG='/path/to/llvm/bin/llvm-config'
```

LLVM support requires a compatible clang compiler (specified, if necessary, using the CLANG environment variable), and a working C++ compiler (specified, if necessary, using the CXX environment variable).
--with-icu

Build with support for the ICU library. This requires the ICU4C package to be installed. The minimum required version of ICU4C is currently 4.2.

By default, pkg-config will be used to find the required compilation options. This is supported for ICU4C version 4.6 and later. For older versions, or if pkg-config is not available, the variables ICU_CFLAGS and ICU_LIBS can be specified to configure, like in this example:

```
./configure ... --with-icu ICU_CFLAGS='-I/some/where/include' ICU_LIBS='-L/some/where/lib -licui18n -licuuc -licudata'
```

(If ICU4C is in the default search path for the compiler, then you still need to specify a nonempty string in order to avoid use of pkg-config, for example, ICU_CFLAGS=' '.)

--with-openssl

Build with support for SSL (encrypted) connections. This requires the OpenSSL package to be installed. configure will check for the required header files and libraries to make sure that your OpenSSL installation is sufficient before proceeding.

--with-pam

Build with PAM (Pluggable Authentication Modules) support.

--with-bsd-auth

Build with BSD Authentication support. (The BSD Authentication framework is currently only available on OpenBSD.)

--with-ldap

Build with LDAP support for authentication and connection parameter lookup (see Section 34.17 and Section 20.10 for more information). On Unix, this requires the OpenLDAP package to be installed. On Windows, the default WinLDAP library is used. configure will check for the required header files and libraries to make sure that your OpenLDAP installation is sufficient before proceeding.

--with-systemd

Build with support for systemd service notifications. This improves integration if the server binary is started under systemd but has no impact otherwise; see Section 18.3 for more information. libsystemd and the associated header files need to be installed to be able to use this option.

--without-readline

Prevents use of the Readline library (and libedit as well). This option disables command-line editing and history in psql, so it is not recommended.

--with-libedit-preferred

Favors the use of the BSD-licensed libedit library rather than GPL-licensed Readline. This option is significant only if you have both libraries installed; the default in that case is to use Readline.
Installation from Source Code

--with-bonjour

Build with Bonjour support. This requires Bonjour support in your operating system. Recommended on macOS.

--with-uuid=LIBRARY

Build the uuid-ossp module (which provides functions to generate UUIDs), using the specified UUID library. LIBRARY must be one of:

- *bsd* to use the UUID functions found in FreeBSD, NetBSD, and some other BSD-derived systems
- *e2fs* to use the UUID library created by the e2fsprogs project; this library is present in most Linux systems and in macOS, and can be obtained for other platforms as well
- *ossp* to use the OSSP UUID library

--with-uuid=oss

Obsolete equivalent of --with-uuid=oss.

--with-libxml

Build with libxml (enables SQL/XML support). Libxml version 2.6.23 or later is required for this feature.

Libxml installs a program `xml2-config` that can be used to detect the required compiler and linker options. PostgreSQL will use it automatically if found. To specify a libxml installation at an unusual location, you can either set the environment variable `XML2_CONFIG` to point to the `xml2-config` program belonging to the installation, or use the options --with-includes and --with-libraries.

--with-libxslt

Use libxslt when building the xml2 module. xml2 relies on this library to perform XSL transformations of XML.

--disable-float4-byval

Disable passing float4 values “by value”, causing them to be passed “by reference” instead. This option costs performance, but may be needed for compatibility with old user-defined functions that are written in C and use the “version 0” calling convention. A better long-term solution is to update any such functions to use the “version 1” calling convention.

--disable-float8-byval

Disable passing float8 values “by value”, causing them to be passed “by reference” instead. This option costs performance, but may be needed for compatibility with old user-defined functions that are written in C and use the “version 0” calling convention. A better long-term solution is to update any such functions to use the “version 1” calling convention. Note that this option affects not only float8, but also int8 and some related types such as timestamp. On 32-bit platforms, --disable-float8-byval is the default and it is not allowed to select --enable-float8-byval.

Installation from Source Code

--with-segsize=SEGSIZE

Set the *segment size*, in gigabytes. Large tables are divided into multiple operating-system files, each of size equal to the segment size. This avoids problems with file size limits that exist on many platforms. The default segment size, 1 gigabyte, is safe on all supported platforms. If your operating system has “largefile” support (which most do, nowadays), you can use a larger segment size. This can be helpful to reduce the number of file descriptors consumed when working with very large tables. But be careful not to select a value larger than is supported by your platform and the file systems you intend to use. Other tools you might wish to use, such as tar, could also set limits on the usable file size. It is recommended, though not absolutely required, that this value be a power of 2. Note that changing this value requires an initdb.

--with-blocksize=BLOCKSIZE

Set the *block size*, in kilobytes. This is the unit of storage and I/O within tables. The default, 8 kilobytes, is suitable for most situations; but other values may be useful in special cases. The value must be a power of 2 between 1 and 32 (kilobytes). Note that changing this value requires an initdb.

--with-wal-blocksize=BLOCKSIZE

Set the *WAL block size*, in kilobytes. This is the unit of storage and I/O within the WAL log. The default, 8 kilobytes, is suitable for most situations; but other values may be useful in special cases. The value must be a power of 2 between 1 and 64 (kilobytes). Note that changing this value requires an initdb.

--disable-spinlocks

Allow the build to succeed even if PostgreSQL has no CPU spinlock support for the platform. The lack of spinlock support will result in poor performance; therefore, this option should only be used if the build aborts and informs you that the platform lacks spinlock support. If this option is required to build PostgreSQL on your platform, please report the problem to the PostgreSQL developers.

--disable-strong-random

Allow the build to succeed even if PostgreSQL has no support for strong random numbers on the platform. A source of random numbers is needed for some authentication protocols, as well as some routines in the pgcrypto module. --disable-strong-random disables functionality that requires cryptographically strong random numbers, and substitutes a weak pseudo-random-number-generator for the generation of authentication salt values and query cancel keys. It may make authentication less secure.

--disable-thread-safety

Disable the thread-safety of client libraries. This prevents concurrent threads in libpq and ECPG programs from safely controlling their private connection handles.

--with-system-tzdata=DIRECTORY

PostgreSQL includes its own time zone database, which it requires for date and time operations. This time zone database is in fact compatible with the IANA time zone database provided by many operating systems such as FreeBSD, Linux, and Solaris, so it would be redundant to install it again. When this option is used, the system-supplied time zone database in DIRECTORY is used instead of the one included in the PostgreSQL source distribution. DIRECTORY must be specified as an absolute path. /usr/share/zoneinfo is a likely directory on some operating systems. Note that the installation routine will not detect mismatching or erroneous time zone data. If you use this option, you are advised to run the
regression tests to verify that the time zone data you have pointed to works correctly with PostgreSQL.

This option is mainly aimed at binary package distributors who know their target operating system well. The main advantage of using this option is that the PostgreSQL package won’t need to be upgraded whenever any of the many local daylight-saving time rules change. Another advantage is that PostgreSQL can be cross-compiled more straightforwardly if the time zone database files do not need to be built during the installation.

--without-zlib

Prevents use of the Zlib library. This disables support for compressed archives in pg_dump and pg_restore. This option is only intended for those rare systems where this library is not available.

--enable-debug

Compiles all programs and libraries with debugging symbols. This means that you can run the programs in a debugger to analyze problems. This enlarges the size of the installed executables considerably, and on non-GCC compilers it usually also disables compiler optimization, causing slowdowns. However, having the symbols available is extremely helpful for dealing with any problems that might arise. Currently, this option is recommended for production installations only if you use GCC. But you should always have it on if you are doing development work or running a beta version.

--enable-coverage

If using GCC, all programs and libraries are compiled with code coverage testing instrumentation. When run, they generate files in the build directory with code coverage metrics. See Section 33.5 for more information. This option is for use only with GCC and when doing development work.

--enable-profiling

If using GCC, all programs and libraries are compiled so they can be profiled. On backend exit, a subdirectory will be created that contains the gmon.out file for use in profiling. This option is for use only with GCC and when doing development work.

--enable-cassert

Enables assertion checks in the server, which test for many “cannot happen” conditions. This is invaluable for code development purposes, but the tests can slow down the server significantly. Also, having the tests turned on won’t necessarily enhance the stability of your server! The assertion checks are not categorized for severity, and so what might be a relatively harmless bug will still lead to server restarts if it triggers an assertion failure. This option is not recommended for production use, but you should have it on for development work or when running a beta version.

--enable-depend

Enables automatic dependency tracking. With this option, the makefiles are set up so that all affected object files will be rebuilt when any header file is changed. This is useful if you are doing development work, but is just wasted overhead if you intend only to compile once and install. At present, this option only works with GCC.

--enable-dtrace

Compiles PostgreSQL with support for the dynamic tracing tool DTrace. See Section 28.5 for more information.
To point to the dtrace program, the environment variable DTRACE can be set. This will often be necessary because dtrace is typically installed under /usr/sbin, which might not be in the path.

Extra command-line options for the dtrace program can be specified in the environment variable DTRACEFLAGS. On Solaris, to include DTrace support in a 64-bit binary, you must specify DTRACEFLAGS="-64" to configure. For example, using the GCC compiler:

```
./configure CC='gcc -m64' --enable-dtrace
DTRACEFLAGS='-64' ...
```

Using Sun's compiler:

```
./configure CC='/opt/SUNWspro/bin/cc -xtarget=native64' --
enable-dtrace DTRACEFLAGS='-64' ...
```

---enable-tap-tests

Enable tests using the Perl TAP tools. This requires a Perl installation and the Perl module IPC::Run. See Section 33.4 for more information.

If you prefer a C compiler different from the one configure picks, you can set the environment variable CC to the program of your choice. By default, configure will pick gcc if available, else the platform's default (usually cc). Similarly, you can override the default compiler flags if needed with the CFLAGS variable.

You can specify environment variables on the configure command line, for example:

```
./configure CC='/opt/bin/gcc CFLAGS='-O2 -pipe'
```

Here is a list of the significant variables that can be set in this manner:

- **BISON**
  
  Bison program

- **CC**
  
  C compiler

- **CFLAGS**
  
  options to pass to the C compiler

- **CLANG**
  
  path to clang program used to process source code for inlining when compiling with --with-llvm

- **CPP**
  
  C preprocessor

- **CPPFLAGS**
  
  options to pass to the C preprocessor
CXX

C++ compiler

CXXFLAGS

options to pass to the C++ compiler

DTRACE

location of the dtrace program

DTRACEFLAGS

options to pass to the dtrace program

FLEX

Flex program

LDFLAGS

options to use when linking either executables or shared libraries

LDFLAGS_EX

additional options for linking executables only

LDFLAGS_SL

additional options for linking shared libraries only

LLVM_CONFIG

llvm-config program used to locate the LLVM installation.

MSGFMT

msgfmt program for native language support

PERL

Full path name of the Perl interpreter. This will be used to determine the dependencies for building PL/Perl.

PYTHON

Full path name of the Python interpreter. This will be used to determine the dependencies for building PL/Python. Also, whether Python 2 or 3 is specified here (or otherwise implicitly chosen) determines which variant of the PL/Python language becomes available. See Section 46.1 for more information.

TCLSH

Full path name of the Tcl interpreter. This will be used to determine the dependencies for building PL/Tcl, and it will be substituted into Tcl scripts.

XML2_CONFIG

xml2-config program used to locate the libxml installation.
Sometimes it is useful to add compiler flags after-the-fact to the set that were chosen by configure. An important example is that gcc's `-Werror` option cannot be included in the CFLAGS passed to configure, because it will break many of configure's built-in tests. To add such flags, include them in the COPT environment variable while running make. The contents of COPT are added to both the CFLAGS and LDFLAGS options set up by configure. For example, you could do

```make
make COPT='-Werror'
```

or

```make
export COPT='-Werror'
make
```

**Note**

When developing code inside the server, it is recommended to use the configure options `--enable-cassert` (which turns on many run-time error checks) and `--enable-debug` (which improves the usefulness of debugging tools).

If using GCC, it is best to build with an optimization level of at least `-O1`, because using no optimization (`-O0`) disables some important compiler warnings (such as the use of uninitialized variables). However, non-zero optimization levels can complicate debugging because stepping through compiled code will usually not match up one-to-one with source code lines. If you get confused while trying to debug optimized code, recompile the specific files of interest with `-O0`. An easy way to do this is by passing an option to make: `make PROFILE=-O0 file.o`.

The COPT and PROFILE environment variables are actually handled identically by the PostgreSQL makefiles. Which to use is a matter of preference, but a common habit among developers is to use PROFILE for one-time flag adjustments, while COPT might be kept set all the time.

2. **Build**

To start the build, type either of:

```make
make
make all
```

(Remember to use GNU make.) The build will take a few minutes depending on your hardware. The last line displayed should be:

```
All of PostgreSQL successfully made. Ready to install.
```

If you want to build everything that can be built, including the documentation (HTML and man pages), and the additional modules (contrib), type instead:

```make
make world
```

The last line displayed should be:
PostgreSQL, contrib, and documentation successfully made. Ready to install.

If you want to invoke the build from another makefile rather than manually, you must unset MAKELEVEL or set it to zero, for instance like this:

```
build-postgresql:
  $(MAKE) -C postgresql MAKELEVEL=0 all
```

Failure to do that can lead to strange error messages, typically about missing header files.

3. **Regression Tests**

If you want to test the newly built server before you install it, you can run the regression tests at this point. The regression tests are a test suite to verify that PostgreSQL runs on your machine in the way the developers expected it to. Type:

```
make check
```

(This won't work as root; do it as an unprivileged user.) See Chapter 33 for detailed information about interpreting the test results. You can repeat this test at any later time by issuing the same command.

4. **Installing the Files**

**Note**

If you are upgrading an existing system be sure to read Section 18.6, which has instructions about upgrading a cluster.

To install PostgreSQL enter:

```
make install
```

This will install files into the directories that were specified in Step 1. Make sure that you have appropriate permissions to write into that area. Normally you need to do this step as root. Alternatively, you can create the target directories in advance and arrange for appropriate permissions to be granted.

To install the documentation (HTML and man pages), enter:

```
make install-docs
```

If you built the world above, type instead:

```
make install-world
```

This also installs the documentation.

You can use `make install-strip` instead of `make install` to strip the executable files and libraries as they are installed. This will save some space. If you built with debugging support, stripping will effectively remove the debugging support, so it should only be done if debugging...
Installation from Source Code

is no longer needed. install-strip tries to do a reasonable job saving space, but it does not have perfect knowledge of how to strip every unneeded byte from an executable file, so if you want to save all the disk space you possibly can, you will have to do manual work.

The standard installation provides all the header files needed for client application development as well as for server-side program development, such as custom functions or data types written in C. (Prior to PostgreSQL 8.0, a separate make install-all-headers command was needed for the latter, but this step has been folded into the standard install.)

Client-only installation: If you want to install only the client applications and interface libraries, then you can use these commands:

```
make -C src/bin install
make -C src/include install
make -C src/interfaces install
make -C doc install
```

src/bin has a few binaries for server-only use, but they are small.

Uninstallation: To undo the installation use the command `make uninstall`. However, this will not remove any created directories.

Cleaning: After the installation you can free disk space by removing the built files from the source tree with the command `make clean`. This will preserve the files made by the configure program, so that you can rebuild everything with `make` later on. To reset the source tree to the state in which it was distributed, use `make distclean`. If you are going to build for several platforms within the same source tree you must do this and re-configure for each platform. (Alternatively, use a separate build tree for each platform, so that the source tree remains unmodified.)

If you perform a build and then discover that your configure options were wrong, or if you change anything that configure investigates (for example, software upgrades), then it’s a good idea to do `make distclean` before reconfiguring and rebuilding. Without this, your changes in configuration choices might not propagate everywhere they need to.

16.5. Post-Installation Setup

16.5.1. Shared Libraries

On some systems with shared libraries you need to tell the system how to find the newly installed shared libraries. The systems on which this is not necessary include FreeBSD, HP-UX, Linux, NetBSD, OpenBSD, and Solaris.

The method to set the shared library search path varies between platforms, but the most widely-used method is to set the environment variable `LD_LIBRARY_PATH` like so: In Bourne shells (sh, ksh, bash, zsh):

```
LD_LIBRARY_PATH=/usr/local/pgsql/lib
export LD_LIBRARY_PATH
```

or in csh or tcsh:

```
setenv LD_LIBRARY_PATH /usr/local/pgsql/lib
```

Replace `/usr/local/pgsql/lib` with whatever you set `--libdir` to in Step 1. You should put these commands into a shell start-up file such as `/etc/profile` or `~/.bash_profile`. 

488
Some good information about the caveats associated with this method can be found at http://xahlee.info/UnixResource_dir/_/ldpath.html.

On some systems it might be preferable to set the environment variable `LD_RUN_PATH` before building.

On Cygwin, put the library directory in the `PATH` or move the `.dll` files into the `bin` directory.

If in doubt, refer to the manual pages of your system (perhaps `ld.so` or `rld`). If you later get a message like:

```
psql: error in loading shared libraries
libpq.so.2.1: cannot open shared object file: No such file or directory
```

then this step was necessary. Simply take care of it then.

If you are on Linux and you have root access, you can run:

```
/sbin/ldconfig /usr/local/pgsql/lib
```

(or equivalent directory) after installation to enable the run-time linker to find the shared libraries faster. Refer to the manual page of `ldconfig` for more information. On FreeBSD, NetBSD, and OpenBSD the command is:

```
/sbin/ldconfig -m /usr/local/pgsql/lib
```

instead. Other systems are not known to have an equivalent command.

### 16.5.2. Environment Variables

If you installed into `/usr/local/pgsql` or some other location that is not searched for programs by default, you should add `/usr/local/pgsql/bin` (or whatever you set `--bindir` to in Step 1) into your `PATH`. Strictly speaking, this is not necessary, but it will make the use of PostgreSQL much more convenient.

To do this, add the following to your shell start-up file, such as `~/.bash_profile` (or `/etc/profile`, if you want it to affect all users):

```
PATH=/usr/local/pgsql/bin:$PATH
export PATH
```

If you are using `csh` or `tcsh`, then use this command:

```
set path = ( /usr/local/pgsql/bin $path )
```

To enable your system to find the man documentation, you need to add lines like the following to a shell start-up file unless you installed into a location that is searched by default:

```
MANPATH=/usr/local/pgsql/share/man:$MANPATH
export MANPATH
```

The environment variables `PGHOST` and `PGPORT` specify to client applications the host and port of the database server, overriding the compiled-in defaults. If you are going to run client applications remotely then it is convenient if every user that plans to use the database sets `PGHOST`. This is not required, however; the settings can be communicated via command line options to most client programs.
16.6. Supported Platforms

A platform (that is, a CPU architecture and operating system combination) is considered supported by the PostgreSQL development community if the code contains provisions to work on that platform and it has recently been verified to build and pass its regression tests on that platform. Currently, most testing of platform compatibility is done automatically by test machines in the PostgreSQL Build Farm\(^2\). If you are interested in using PostgreSQL on a platform that is not represented in the build farm, but on which the code works or can be made to work, you are strongly encouraged to set up a build farm member machine so that continued compatibility can be assured.

In general, PostgreSQL can be expected to work on these CPU architectures: x86, x86_64, IA64, PowerPC, PowerPC 64, S/390, S/390x, Sparc, Sparc 64, ARM, MIPS, MIPSEL, and PA-RISC. Code support exists for M68K, M32R, and VAX, but these architectures are not known to have been tested recently. It is often possible to build on an unsupported CPU type by configuring with \texttt{--disable-spinlocks}, but performance will be poor.

PostgreSQL can be expected to work on these operating systems: Linux (all recent distributions), Windows (Win2000 SP4 and later), FreeBSD, OpenBSD, NetBSD, macOS, AIX, HP/UX, and Solaris. Other Unix-like systems may also work but are not currently being tested. In most cases, all CPU architectures supported by a given operating system will work. Look in Section 16.7 below to see if there is information specific to your operating system, particularly if using an older system.

If you have installation problems on a platform that is known to be supported according to recent build farm results, please report it to \texttt{pgsql-bugs@postgresql.org}. If you are interested in porting PostgreSQL to a new platform, \texttt{pgsql-hackers@postgresql.org} is the appropriate place to discuss that.

16.7. Platform-specific Notes

This section documents additional platform-specific issues regarding the installation and setup of PostgreSQL. Be sure to read the installation instructions, and in particular Section 16.2 as well. Also, check Chapter 33 regarding the interpretation of regression test results.

Platforms that are not covered here have no known platform-specific installation issues.

16.7.1. AIX

PostgreSQL works on AIX, but getting it installed properly can be challenging. AIX versions from 4.3.3 to 6.1 are considered supported. You can use GCC or the native IBM compiler \texttt{xlc}. In general, using recent versions of AIX and PostgreSQL helps. Check the build farm for up to date information about which versions of AIX are known to work.

The minimum recommended fix levels for supported AIX versions are:

AIX 4.3.3

Maintenance Level 11 + post ML11 bundle

AIX 5.1

Maintenance Level 9 + post ML9 bundle

AIX 5.2

Technology Level 10 Service Pack 3

\(^2\) https://buildfarm.postgresql.org/
Installation from Source Code

AIX 5.3
Technology Level 7

AIX 6.1
Base Level

To check your current fix level, use oslevel -r in AIX 4.3.3 to AIX 5.2 ML 7, or oslevel -s in later versions.

Use the following configure flags in addition to your own if you have installed Readline or libz in /usr/local: --with-includes=/usr/local/include --with-libraries=/usr/local/lib.

16.7.1.1. GCC Issues

On AIX 5.3, there have been some problems getting PostgreSQL to compile and run using GCC.

You will want to use a version of GCC subsequent to 3.3.2, particularly if you use a prepackaged version. We had good success with 4.0.1. Problems with earlier versions seem to have more to do with the way IBM packaged GCC than with actual issues with GCC, so that if you compile GCC yourself, you might well have success with an earlier version of GCC.

16.7.1.2. Unix-Domain Sockets Broken

AIX 5.3 has a problem where sockaddr_storage is not defined to be large enough. In version 5.3, IBM increased the size of sockaddr_un, the address structure for Unix-domain sockets, but did not correspondingly increase the size of sockaddr_storage. The result of this is that attempts to use Unix-domain sockets with PostgreSQL lead to libpq overflowing the data structure. TCP/IP connections work OK, but not Unix-domain sockets, which prevents the regression tests from working.

The problem was reported to IBM, and is recorded as bug report PMR29657. If you upgrade to maintenance level 5300-03 or later, that will include this fix. A quick workaround is to alter _SS_MAXSIZE to 1025 in /usr/include/sys/socket.h. In either case, recompile PostgreSQL once you have the corrected header file.

16.7.1.3. Internet Address Issues

PostgreSQL relies on the system's getaddrinfo function to parse IP addresses in listen_addresses, pg_hba.conf, etc. Older versions of AIX have assorted bugs in this function. If you have problems related to these settings, updating to the appropriate AIX fix level shown above should take care of it.

One user reports:

When implementing PostgreSQL version 8.1 on AIX 5.3, we periodically ran into problems where the statistics collector would “mysteriously” not come up successfully. This appears to be the result of unexpected behavior in the IPv6 implementation. It looks like PostgreSQL and IPv6 do not play very well together on AIX 5.3.

Any of the following actions “fix” the problem.

• Delete the IPv6 address for localhost:

    (as root)
    # ifconfig lo0 inet6 ::1/0 delete

• Remove IPv6 from net services. The file /etc/netsvc.conf on AIX is roughly equivalent to /etc/nsswitch.conf on Solaris/Linux. The default, on AIX, is thus:
Installation from Source Code

hosts=local,bind

Replace this with:

hosts=local4,bind4

to deactivate searching for IPv6 addresses.

**Warning**

This is really a workaround for problems relating to immaturity of IPv6 support, which improved visibly during the course of AIX 5.3 releases. It has worked with AIX version 5.3, but does not represent an elegant solution to the problem. It has been reported that this workaround is not only unnecessary, but causes problems on AIX 6.1, where IPv6 support has become more mature.

16.7.1.4. Memory Management

AIX can be somewhat peculiar with regards to the way it does memory management. You can have a server with many multiples of gigabytes of RAM free, but still get out of memory or address space errors when running applications. One example is loading of extensions failing with unusual errors. For example, running as the owner of the PostgreSQL installation:

```bash
=# CREATE EXTENSION plperl;
ERROR:  could not load library "/opt/dbs/pgsql/lib/plperl.so": A memory address is not in the address space for the process.
```

Running as a non-owner in the group possessing the PostgreSQL installation:

```bash
=# CREATE EXTENSION plperl;
ERROR:  could not load library "/opt/dbs/pgsql/lib/plperl.so": Bad address
```

Another example is out of memory errors in the PostgreSQL server logs, with every memory allocation near or greater than 256 MB failing.

The overall cause of all these problems is the default bittedness and memory model used by the server process. By default, all binaries built on AIX are 32-bit. This does not depend upon hardware type or kernel in use. These 32-bit processes are limited to 4 GB of memory laid out in 256 MB segments using one of a few models. The default allows for less than 256 MB in the heap as it shares a single segment with the stack.

In the case of the `plperl` example, above, check your umask and the permissions of the binaries in your PostgreSQL installation. The binaries involved in that example were 32-bit and installed as mode 750 instead of 755. Due to the permissions being set in this fashion, only the owner or a member of the possessing group can load the library. Since it isn’t world-readable, the loader places the object into the process’ heap instead of the shared library segments where it would otherwise be placed.

The “ideal” solution for this is to use a 64-bit build of PostgreSQL, but that is not always practical, because systems with 32-bit processors can build, but not run, 64-bit binaries.

If a 32-bit binary is desired, set LDR_CNTRL to MAXDATA=0x{n}00000000, where 1 <= n <= 8, before starting the PostgreSQL server, and try different values and `postgresql.conf` settings to find a
configuration that works satisfactorily. This use of LDR_CNTRL tells AIX that you want the server to have MAXDATA bytes set aside for the heap, allocated in 256 MB segments. When you find a workable configuration, ldedit can be used to modify the binaries so that they default to using the desired heap size. PostgreSQL can also be rebuilt, passing configure LDFLAGS="-Wl,-bmaxdata:0x00000000" to achieve the same effect.

For a 64-bit build, set OBJECT_MODE to 64 and pass CC="gcc -maix64" and LDFLAGS="-Wl,-bbigtoc" to configure. (Options for xlc might differ.) If you omit the export of OBJECT_MODE, your build may fail with linker errors. When OBJECT_MODE is set, it tells AIX's build utilities such as ar, as, and ld what type of objects to default to handling.

By default, overcommit of paging space can happen. While we have not seen this occur, AIX will kill processes when it runs out of memory and the overcommit is accessed. The closest to this that we have seen is fork failing because the system decided that there was not enough memory for another process. Like many other parts of AIX, the paging space allocation method and out-of-memory kill is configurable on a system- or process-wide basis if this becomes a problem.

References and Resources

“Large Program Support1”. AIX Documentation: General Programming Concepts: Writing and Debugging Programs.

“Program Address Space Overview2”. AIX Documentation: General Programming Concepts: Writing and Debugging Programs.

“Performance Overview of the Virtual Memory Manager (VMM)3”. AIX Documentation: Performance Management Guide.


Developing and Porting C and C++ Applications on AIX6. IBM Redbook.

16.7.2. Cygwin

PostgreSQL can be built using Cygwin, a Linux-like environment for Windows, but that method is inferior to the native Windows build (see Chapter 17) and running a server under Cygwin is no longer recommended.

When building from source, proceed according to the normal installation procedure (i.e., ./configure; make; etc.), noting the following-Cygwin specific differences:

• Set your path to use the Cygwin bin directory before the Windows utilities. This will help prevent problems with compilation.

• The adduser command is not supported; use the appropriate user management application on Windows NT, 2000, or XP. Otherwise, skip this step.

• The su command is not supported; use ssh to simulate su on Windows NT, 2000, or XP. Otherwise, skip this step.

• OpenSSL is not supported.

6 http://www.redbooks.ibm.com/abstracts/sg245674.html?Open
• Start cygserver for shared memory support. To do this, enter the command `/usr/sbin/cygserver &`. This program needs to be running anytime you start the PostgreSQL server or initialize a database cluster (`initdb`). The default cygserver configuration may need to be changed (e.g., increase `SEMMNS`) to prevent PostgreSQL from failing due to a lack of system resources.

• Building might fail on some systems where a locale other than C is in use. To fix this, set the locale to C by doing `export LANG=C.utf8` before building, and then setting it back to the previous setting, after you have installed PostgreSQL.

• The parallel regression tests (`make check`) can generate spurious regression test failures due to overflowing the `listen()` backlog queue which causes connection refused errors or hangs. You can limit the number of connections using the make variable `MAX_CONNECTIONS` thus:

  `make MAX_CONNECTIONS=5 check`

  (On some systems you can have up to about 10 simultaneous connections).

It is possible to install cygserver and the PostgreSQL server as Windows NT services. For information on how to do this, please refer to the README document included with the PostgreSQL binary package on Cygwin. It is installed in the directory `/usr/share/doc/Cygwin`.

### 16.7.3. HP-UX

PostgreSQL 7.3+ should work on Series 700/800 PA-RISC machines running HP-UX 10.X or 11.X, given appropriate system patch levels and build tools. At least one developer routinely tests on HP-UX 10.20, and we have reports of successful installations on HP-UX 11.00 and 11.11.

Aside from the PostgreSQL source distribution, you will need GNU make (HP's make will not do), and either GCC or HP's full ANSI C compiler. If you intend to build from Git sources rather than a distribution tarball, you will also need Flex (GNU lex) and Bison (GNU yacc). We also recommend making sure you are fairly up-to-date on HP patches. At a minimum, if you are building 64 bit binaries on HP-UX 11.11 you may need PHSS_30966 (11.11) or a successor patch otherwise `initdb` may hang:

- PHSS_30966  s700_800 ld(1) and linker tools cumulative patch

  On general principles you should be current on libc and ld/dld patches, as well as compiler patches if you are using HP's C compiler. See HP's support sites such as ftp://us-ffs.external.hp.com/ for free copies of their latest patches.

  If you are building on a PA-RISC 2.0 machine and want to have 64-bit binaries using GCC, you must use a GCC 64-bit version.

  If you are building on a PA-RISC 2.0 machine and want the compiled binaries to run on PA-RISC 1.1 machines you will need to specify `+DApportable` in `CFLAGS`.

  If you are building on a HP-UX Itanium machine, you will need the latest HP ANSI C compiler with its dependent patch or successor patches:

- PHSS_30848  s700_800 HP C Compiler (A.05.57)
- PHSS_30849  s700_800 u2comp/be/plugin library Patch

If you have both HP's C compiler and GCC's, then you might want to explicitly select the compiler to use when you run `configure`:  

---

494
Installation from Source Code

./configure CC=cc

for HP's C compiler, or

./configure CC=gcc

for GCC. If you omit this setting, then configure will pick gcc if it has a choice.

The default install target location is /usr/local/pgsql, which you might want to change to something under /opt. If so, use the --prefix switch to configure.

In the regression tests, there might be some low-order-digit differences in the geometry tests, which vary depending on which compiler and math library versions you use. Any other error is cause for suspicion.

16.7.4. macOS

On recent macOS releases, it’s necessary to embed the “sysroot” path in the include switches used to find some system header files. This results in the outputs of the configure script varying depending on which SDK version was used during configure. That shouldn’t pose any problem in simple scenarios, but if you are trying to do something like building an extension on a different machine than the server code was built on, you may need to force use of a different sysroot path. To do that, set PG_SYSROOT, for example

make PG_SYSROOT=/desired/path all

To find out the appropriate path on your machine, run

xcodebuild -version -sdk macosx Path

Note that building an extension using a different sysroot version than was used to build the core server is not really recommended; in the worst case it could result in hard-to-debug ABI inconsistencies.

You can also select a non-default sysroot path when configuring, by specifying PG_SYSROOT to configure:

./configure ... PG_SYSROOT=/desired/path

macOS’s “System Integrity Protection” (SIP) feature breaks make check, because it prevents passing the needed setting of DYLD_LIBRARY_PATH down to the executables being tested. You can work around that by doing make install before make check. Most Postgres developers just turn off SIP, though.

16.7.5. MinGW/Native Windows

PostgreSQL for Windows can be built using MinGW, a Unix-like build environment for Microsoft operating systems, or using Microsoft's Visual C++ compiler suite. The MinGW build variant uses the normal build system described in this chapter; the Visual C++ build works completely differently and is described in Chapter 17. It is a fully native build and uses no additional software like MinGW. A ready-made installer is available on the main PostgreSQL web site.

The native Windows port requires a 32 or 64-bit version of Windows 2000 or later. Earlier operating systems do not have sufficient infrastructure (but Cygwin may be used on those). MinGW, the Unix-like build tools, and MSYS, a collection of Unix tools required to run shell scripts like configure, can be downloaded from http://www.mingw.org/. Neither is required to run the resulting binaries; they are needed only for creating the binaries.
To build 64 bit binaries using MinGW, install the 64 bit tool set from https://mingw-w64.org/, put its bin directory in the PATH, and run configure with the --host=x86_64-w64-mingw32 option.

After you have everything installed, it is suggested that you run psql under CMD.EXE, as the MSYS console has buffering issues.

16.7.5.1. Collecting Crash Dumps on Windows

If PostgreSQL on Windows crashes, it has the ability to generate minidumps that can be used to track down the cause for the crash, similar to core dumps on Unix. These dumps can be read using the Windows Debugger Tools or using Visual Studio. To enable the generation of dumps on Windows, create a subdirectory named crashdumps inside the cluster data directory. The dumps will then be written into this directory with a unique name based on the identifier of the crashing process and the current time of the crash.

16.7.6. Solaris

PostgreSQL is well-supported on Solaris. The more up to date your operating system, the fewer issues you will experience; details below.

16.7.6.1. Required Tools

You can build with either GCC or Sun's compiler suite. For better code optimization, Sun's compiler is strongly recommended on the SPARC architecture. We have heard reports of problems when using GCC 2.95.1; GCC 2.95.3 or later is recommended. If you are using Sun's compiler, be careful not to select /usr/ucb/cc; use /opt/SUNWspro/bin/cc.

You can download Sun Studio from https://www.oracle.com/technetwork/server-storage/solarisstudio/downloads/ . Many of GNU tools are integrated into Solaris 10, or they are present on the Solaris companion CD. If you like packages for older version of Solaris, you can find these tools at http://www.sunfreeware.com . If you prefer sources, look at https://www.gnu.org/prep/ftp.

16.7.6.2. configure Complains About a Failed Test Program

If configure complains about a failed test program, this is probably a case of the run-time linker being unable to find some library, probably libz, libreadline or some other non-standard library such as libssl. To point it to the right location, set the LDFLAGS environment variable on the configure command line, e.g.,

```
configure ...
```

LDFLAGS="-R /usr/sfw/lib:/opt/sfw/lib:/usr/local/lib"

See the ld man page for more information.

16.7.6.3. 64-bit Build Sometimes Crashes

On Solaris 7 and older, the 64-bit version of libc has a buggy vsnprintf routine, which leads to erratic core dumps in PostgreSQL. The simplest known workaround is to force PostgreSQL to use its own version of vsnprintf rather than the library copy. To do this, after you run configure edit a file produced by configure: In src/Makefile.global, change the line

```
LIBOBJS =
```

to read

```
LIBOBJS = snprintf.o
```
Installation from Source Code

(There might be other files already listed in this variable. Order does not matter.) Then build as usual.

16.7.6.4. Compiling for Optimal Performance

On the SPARC architecture, Sun Studio is strongly recommended for compilation. Try using the `-xO5` optimization flag to generate significantly faster binaries. Do not use any flags that modify behavior of floating-point operations and `errno` processing (e.g., `-fast`). These flags could raise some nonstandard PostgreSQL behavior for example in the date/time computing.

If you do not have a reason to use 64-bit binaries on SPARC, prefer the 32-bit version. The 64-bit operations are slower and 64-bit binaries are slower than the 32-bit variants. And on other hand, 32-bit code on the AMD64 CPU family is not native, and that is why 32-bit code is significant slower on this CPU family.

16.7.6.5. Using DTrace for Tracing PostgreSQL

Yes, using DTrace is possible. See Section 28.5 for further information.

If you see the linking of the `postgres` executable abort with an error message like:

```
Undefined                       first referenced
 symbol                        in file
AbortTransaction               utils/probes.o
CommitTransaction              utils/probes.o
ld: fatal: Symbol referencing errors. No output written to postgres
collect2: ld returned 1 exit status
make: *** [postgres] Error 1
```

your DTrace installation is too old to handle probes in static functions. You need Solaris 10u4 or newer.
Chapter 17. Installation from Source Code on Windows

It is recommended that most users download the binary distribution for Windows, available as a graphical installer package from the PostgreSQL website. Building from source is only intended for people developing PostgreSQL or extensions.

There are several different ways of building PostgreSQL on Windows. The simplest way to build with Microsoft tools is to install Visual Studio Express 2017 for Windows Desktop and use the included compiler. It is also possible to build with the full Microsoft Visual C++ 2005 to 2017. In some cases that requires the installation of the Windows SDK in addition to the compiler.

It is also possible to build PostgreSQL using the GNU compiler tools provided by MinGW, or using Cygwin for older versions of Windows.

Building using MinGW or Cygwin uses the normal build system, see Chapter 16 and the specific notes in Section 16.7.5 and Section 16.7.2. To produce native 64 bit binaries in these environments, use the tools from MinGW-w64. These tools can also be used to cross-compile for 32 bit and 64 bit Windows targets on other hosts, such as Linux and macOS. Cygwin is not recommended for running a production server, and it should only be used for running on older versions of Windows where the native build does not work, such as Windows 98. The official binaries are built using Visual Studio.

Native builds of psql don't support command line editing. The Cygwin build does support command line editing, so it should be used where psql is needed for interactive use on Windows.

17.1. Building with Visual C++ or the Microsoft Windows SDK

PostgreSQL can be built using the Visual C++ compiler suite from Microsoft. These compilers can be either from Visual Studio, Visual Studio Express or some versions of the Microsoft Windows SDK. If you do not already have a Visual Studio environment set up, the easiest ways are to use the compilers from Visual Studio Express 2017 for Windows Desktop or those in the Windows SDK 8.1, which are both free downloads from Microsoft.

Both 32-bit and 64-bit builds are possible with the Microsoft Compiler suite. 32-bit PostgreSQL builds are possible with Visual Studio 2005 to Visual Studio 2017 (including Express editions), as well as standalone Windows SDK releases 6.0 to 8.1. 64-bit PostgreSQL builds are supported with Microsoft Windows SDK version 6.0a to 8.1 or Visual Studio 2008 and above. Compilation is supported down to Windows XP and Windows Server 2003 when building with Visual Studio 2005 to Visual Studio 2013. Building with Visual Studio 2015 is supported down to Windows Vista and Windows Server 2008. Building with Visual Studio 2017 is supported down to Windows 7 SP1 and Windows Server 2008 R2 SP1.

The tools for building using Visual C++ or Platform SDK are in the src/tools/msvc directory. When building, make sure there are no tools from MinGW or Cygwin present in your system PATH. Also, make sure you have all the required Visual C++ tools available in the PATH. In Visual Studio, start the Visual Studio Command Prompt. If you wish to build a 64-bit version, you must use the 64-bit version of the command, and vice versa. In the Microsoft Windows SDK, start the CMD shell listed under the SDK on the Start Menu. In recent SDK versions you can change the targeted CPU architecture, build type, and target OS by using the setenv command, e.g. setenv /x86 /release /xp to target Windows XP or later with a 32-bit release build. See /? for other options to setenv. All commands should be run from the src\tools\msvc directory.

Before you build, you may need to edit the file config.pl to reflect any configuration options you want to change, or the paths to any third party libraries to use. The complete configuration is
determined by first reading and parsing the file `config_default.pl`, and then apply any changes from `config.pl`. For example, to specify the location of your Python installation, put the following in `config.pl`:

```perl
$config->{python} = 'c:\python26';
```

You only need to specify those parameters that are different from what's in `config_default.pl`. If you need to set any other environment variables, create a file called `buildenv.pl` and put the required commands there. For example, to add the path for bison when it's not in the PATH, create a file containing:

```perl
$ENV{PATH}=$ENV{PATH} . ';c:\some\where\bison\bin';
```

To pass additional command line arguments to the Visual Studio build command (msbuild or vcbuild):

```perl
$ENV{MSBFLAGS}="/m";
```

### 17.1.1. Requirements

The following additional products are required to build PostgreSQL. Use the `config.pl` file to specify which directories the libraries are available in.

**Microsoft Windows SDK**

If your build environment doesn't ship with a supported version of the Microsoft Windows SDK it is recommended that you upgrade to the latest version (currently version 7.1), available for download from https://www.microsoft.com/download.

You must always include the Windows Headers and Libraries part of the SDK. If you install a Windows SDK including the Visual C++ Compilers, you don't need Visual Studio to build. Note that as of Version 8.0a the Windows SDK no longer ships with a complete command-line build environment.

**ActiveState Perl**

ActiveState Perl is required to run the build generation scripts. MinGW or Cygwin Perl will not work. It must also be present in the PATH. Binaries can be downloaded from https://www.activestate.com (Note: version 5.8.3 or later is required, the free Standard Distribution is sufficient).

The following additional products are not required to get started, but are required to build the complete package. Use the `config.pl` file to specify which directories the libraries are available in.

**ActiveState TCL**

Required for building PL/Tcl (Note: version 8.4 is required, the free Standard Distribution is sufficient).

**Bison and Flex**

Bison and Flex are required to build from Git, but not required when building from a release file. Only Bison 1.875 or versions 2.2 and later will work. Flex must be version 2.5.31 or later.

Both Bison and Flex are included in the msys tool suite, available from http://www.mingw.org/wiki/MSYS as part of the MinGW compiler suite.

You will need to add the directory containing `flex.exe` and `bison.exe` to the PATH environment variable in `buildenv.pl` unless they are already in PATH. In the case of MinGW, the directory is the `\msys\1.0\bin` subdirectory of your MinGW installation directory.
Installation from Source
Code on Windows

Note

The Bison distribution from GnuWin32 appears to have a bug that causes Bison to malfunction when installed in a directory with spaces in the name, such as the default location on English installations C:\Program Files \GnuWin32. Consider installing into C:\GnuWin32 or use the NTFS short name path to GnuWin32 in your PATH environment setting (e.g. C:\PROGRA~1\GnuWin32).

Note

The obsolete winflex binaries distributed on the PostgreSQL FTP site and referenced in older documentation will fail with “flex: fatal internal error, exec failed” on 64-bit Windows hosts. Use Flex from MSYS instead.

Diff

Diff is required to run the regression tests, and can be downloaded from http://gnuwin32.sourceforge.net.

Gettext

Gettext is required to build with NLS support, and can be downloaded from http://gnuwin32.sourceforge.net. Note that binaries, dependencies and developer files are all needed.

MIT Kerberos

Required for GSSAPI authentication support. MIT Kerberos can be downloaded from http://web.mit.edu/Kerberos/dist/index.html.

libxml2 and libxslt

Required for XML support. Binaries can be downloaded from http://zlatkovic.com/pub/libxml or source from http://xmlsoft.org. Note that libxml2 requires iconv, which is available from the same download location.

OpenSSL


ossp-uuid

Required for UUID-OSSP support (contrib only). Source can be downloaded from http://www.ossp.org/pkg/lib/uuid/.

Python

Required for building PL/Python. Binaries can be downloaded from https://www.python.org.

zlib

Required for compression support in pg_dump and pg_restore. Binaries can be downloaded from http://www.zlib.net.
17.1.2. Special Considerations for 64-bit Windows

PostgreSQL will only build for the x64 architecture on 64-bit Windows, there is no support for Itanium processors.

Mixing 32- and 64-bit versions in the same build tree is not supported. The build system will automatically detect if it's running in a 32- or 64-bit environment, and build PostgreSQL accordingly. For this reason, it is important to start the correct command prompt before building.

To use a server-side third party library such as python or OpenSSL, this library must also be 64-bit. There is no support for loading a 32-bit library in a 64-bit server. Several of the third party libraries that PostgreSQL supports may only be available in 32-bit versions, in which case they cannot be used with 64-bit PostgreSQL.

17.1.3. Building

To build all of PostgreSQL in release configuration (the default), run the command:

`build`

To build all of PostgreSQL in debug configuration, run the command:

`build DEBUG`

To build just a single project, for example psql, run the commands:

`build psql`
`build DEBUG psql`

To change the default build configuration to debug, put the following in the `buildenv.pl` file:

```bash
$ENV{CONFIG}="Debug";
```

It is also possible to build from inside the Visual Studio GUI. In this case, you need to run:

`perl mkvcbuild.pl`

from the command prompt, and then open the generated `pgsql.sln` (in the root directory of the source tree) in Visual Studio.

17.1.4. Cleaning and Installing

Most of the time, the automatic dependency tracking in Visual Studio will handle changed files. But if there have been large changes, you may need to clean the installation. To do this, simply run the `clean.bat` command, which will automatically clean out all generated files. You can also run it with the `dist` parameter, in which case it will behave like `make distclean` and remove the flex/bison output files as well.

By default, all files are written into a subdirectory of the debug or release directories. To install these files using the standard layout, and also generate the files required to initialize and use the database, run the command:

`install c:\destination\directory`
If you want to install only the client applications and interface libraries, then you can use these commands:

```
install c:\destination\directory client
```

### 17.1.5. Running the Regression Tests

To run the regression tests, make sure you have completed the build of all required parts first. Also, make sure that the DLLs required to load all parts of the system (such as the Perl and Python DLLs for the procedural languages) are present in the system path. If they are not, set it through the `buildenv.pl` file. To run the tests, run one of the following commands from the `src\tools \msvc` directory:

```
vcregress check
vcregress installcheck
vcregress plcheck
vcregress contribcheck
vcregress modulescheck
vcregress ecpgcheck
vcregress isolationcheck
vcregress bincheck
vcregress recoverycheck
vcregress upgradecheck
```

To change the schedule used (default is parallel), append it to the command line like:

```
vcregress check serial
```

For more information about the regression tests, see Chapter 33.

Running the regression tests on client programs, with `vcregress bincheck`, or on recovery tests, with `vcregress recoverycheck`, requires an additional Perl module to be installed:

**IPC::Run**

As of this writing, `IPC::Run` is not included in the ActiveState Perl installation, nor in the ActiveState Perl Package Manager (PPM) library. To install, download the `IPC-Run-<version>.tar.gz` source archive from CPAN, at https://metacpan.org/release/IPC-Run, and uncompress. Edit the `buildenv.pl` file, and add a PERL5LIB variable to point to the `lib` subdirectory from the extracted archive. For example:

```
$ENV{PERL5LIB}=$ENV{PERL5LIB} . ';c:\IPC-Run-0.94\lib';
```

### 17.1.6. Building the Documentation

Building the PostgreSQL documentation in HTML format requires several tools and files. Create a root directory for all these files, and store them in the subdirectories in the list below.

**OpenJade 1.3.1-2**

Download from https://sourceforge.net/projects/openjade/files/openjade/1.3.1/openjade-1.3.1-2-bin.zip/download and uncompress in the subdirectory `openjade-1.3.1`.

**DocBook DTD 4.2**

Download from https://www.oasis-open.org/docbook/sgml/4.2/docbook-4.2.zip and uncompress in the subdirectory `docbook`.

502
ISO character entities

Download from https://www.oasis-open.org/cover/ISOEnts.zip and uncompressed in the subdirectory docbook.

Edit the buildenv.pl file, and add a variable for the location of the root directory, for example:

```bash
$ENV{DOCROOT}='c:\docbook';
```

To build the documentation, run the command builddoc.bat. Note that this will actually run the build twice, in order to generate the indexes. The generated HTML files will be in `doc\src\sgml`. 
Chapter 18. Server Setup and Operation

This chapter discusses how to set up and run the database server and its interactions with the operating system.

18.1. The PostgreSQL User Account

As with any server daemon that is accessible to the outside world, it is advisable to run PostgreSQL under a separate user account. This user account should only own the data that is managed by the server, and should not be shared with other daemons. (For example, using the user nobody is a bad idea.) It is not advisable to install executables owned by this user because compromised systems could then modify their own binaries.

To add a Unix user account to your system, look for a command useradd or adduser. The user name postgres is often used, and is assumed throughout this book, but you can use another name if you like.

18.2. Creating a Database Cluster

Before you can do anything, you must initialize a database storage area on disk. We call this a database cluster. (The SQL standard uses the term catalog cluster.) A database cluster is a collection of databases that is managed by a single instance of a running database server. After initialization, a database cluster will contain a database named postgres, which is meant as a default database for use by utilities, users and third party applications. The database server itself does not require the postgres database to exist, but many external utility programs assume it exists. Another database created within each cluster during initialization is called template1. As the name suggests, this will be used as a template for subsequently created databases; it should not be used for actual work. (See Chapter 22 for information about creating new databases within a cluster.)

In file system terms, a database cluster is a single directory under which all data will be stored. We call this the data directory or data area. It is completely up to you where you choose to store your data. There is no default, although locations such as /usr/local/pgsql/data or /var/lib/pgsql/data are popular. To initialize a database cluster, use the command initdb, which is installed with PostgreSQL. The desired file system location of your database cluster is indicated by the -D option, for example:

```
$ initdb -D /usr/local/pgsql/data
```

Note that you must execute this command while logged into the PostgreSQL user account, which is described in the previous section.

Tip

As an alternative to the -D option, you can set the environment variable PGDATA.

Alternatively, you can run initdb via the pg_ctl program like so:

```
$ pg_ctl -D /usr/local/pgsql/data initdb
```
This may be more intuitive if you are using `pg_ctl` for starting and stopping the server (see Section 18.3), so that `pg_ctl` would be the sole command you use for managing the database server instance.

`initdb` will attempt to create the directory you specify if it does not already exist. Of course, this will fail if `initdb` does not have permissions to write in the parent directory. It's generally recommendable that the PostgreSQL user own not just the data directory but its parent directory as well, so that this should not be a problem. If the desired parent directory doesn't exist either, you will need to create it first, using root privileges if the grandparent directory isn't writable. So the process might look like this:

```
root# mkdir /usr/local/pgsql
root# chown postgres /usr/local/pgsql
root# su postgres
postgres$ initdb -D /usr/local/pgsql/data
```

`initdb` will refuse to run if the data directory exists and already contains files; this is to prevent accidentally overwriting an existing installation.

Because the data directory contains all the data stored in the database, it is essential that it be secured from unauthorized access. `initdb` therefore revokes access permissions from everyone but the PostgreSQL user, and optionally, group. Group access, when enabled, is read-only. This allows an unprivileged user in the same group as the cluster owner to take a backup of the cluster data or perform other operations that only require read access.

Note that enabling or disabling group access on an existing cluster requires the cluster to be shut down and the appropriate mode to be set on all directories and files before restarting PostgreSQL. Otherwise, a mix of modes might exist in the data directory. For clusters that allow access only by the owner, the appropriate modes are `0700` for directories and `0600` for files. For clusters that also allow reads by the group, the appropriate modes are `0750` for directories and `0640` for files.

However, while the directory contents are secure, the default client authentication setup allows any local user to connect to the database and even become the database superuser. If you do not trust other local users, we recommend you use one of `initdb`'s `-W`, `--pwprompt` or `--pwfile` options to assign a password to the database superuser. Also, specify `-A md5` or `-A password` so that the default trust authentication mode is not used; or modify the generated `pg_hba.conf` file after running `initdb`, but before you start the server for the first time. (Other reasonable approaches include using peer authentication or file system permissions to restrict connections. See Chapter 20 for more information.)

`initdb` also initializes the default locale for the database cluster. Normally, it will just take the locale settings in the environment and apply them to the initialized database. It is possible to specify a different locale for the database; more information about that can be found in Section 23.1. The default sort order used within the particular database cluster is set by `initdb`, and while you can create new databases using different sort order, the order used in the template databases that `initdb` creates cannot be changed without dropping and recreating them. There is also a performance impact for using locales other than C or POSIX. Therefore, it is important to make this choice correctly the first time.

`initdb` also sets the default character set encoding for the database cluster. Normally this should be chosen to match the locale setting. For details see Section 23.3.

Non-C and non-POSIX locales rely on the operating system's collation library for character set ordering. This controls the ordering of keys stored in indexes. For this reason, a cluster cannot switch to an incompatible collation library version, either through snapshot restore, binary streaming replication, a different operating system, or an operating system upgrade.

### 18.2.1. Use of Secondary File Systems

Many installations create their database clusters on file systems (volumes) other than the machine's "root" volume. If you choose to do this, it is not advisable to try to use the secondary volume's topmost
directory (mount point) as the data directory. Best practice is to create a directory within the mount-point directory that is owned by the PostgreSQL user, and then create the data directory within that. This avoids permissions problems, particularly for operations such as pg_upgrade, and it also ensures clean failures if the secondary volume is taken offline.

18.2.2. Use of Network File Systems

Many installations create their database clusters on network file systems. Sometimes this is done via NFS, or by using a Network Attached Storage (NAS) device that uses NFS internally. PostgreSQL does nothing special for NFS file systems, meaning it assumes NFS behaves exactly like locally-connected drives. If the client or server NFS implementation does not provide standard file system semantics, this can cause reliability problems (see https://www.time-travellers.org/shane/papers/NFS_considered_harmful.html). Specifically, delayed (asynchronous) writes to the NFS server can cause data corruption problems. If possible, mount the NFS file system synchronously (without caching) to avoid this hazard. Also, soft-mounting the NFS file system is not recommended.

Storage Area Networks (SAN) typically use communication protocols other than NFS, and may or may not be subject to hazards of this sort. It's advisable to consult the vendor's documentation concerning data consistency guarantees. PostgreSQL cannot be more reliable than the file system it's using.

18.3. Starting the Database Server

Before anyone can access the database, you must start the database server. The database server program is called postgres. The postgres program must know where to find the data it is supposed to use. This is done with the -D option. Thus, the simplest way to start the server is:

```
$ postgres -D /usr/local/pgsql/data
```

which will leave the server running in the foreground. This must be done while logged into the PostgreSQL user account. Without -D, the server will try to use the data directory named by the environment variable PGDATA. If that variable is not provided either, it will fail.

Normally it is better to start postgres in the background. For this, use the usual Unix shell syntax:

```
$ postgres -D /usr/local/pgsql/data >logfile 2>&1 &
```

It is important to store the server's stdout and stderr output somewhere, as shown above. It will help for auditing purposes and to diagnose problems. (See Section 24.3 for a more thorough discussion of log file handling.)

The postgres program also takes a number of other command-line options. For more information, see the postgres reference page and Chapter 19 below.

This shell syntax can get tedious quickly. Therefore the wrapper program pg_ctl is provided to simplify some tasks. For example:

```
pg_ctl start -l logfile
```

will start the server in the background and put the output into the named log file. The -D option has the same meaning here as for postgres. pg_ctl is also capable of stopping the server.

Normally, you will want to start the database server when the computer boots. Autostart scripts are operating-system-specific. There are a few distributed with PostgreSQL in the contrib/start-scripts directory. Installing one will require root privileges.

Different systems have different conventions for starting up daemons at boot time. Many systems have a file /etc/rc.local or /etc/rc.d/rc.local. Others use init.d or rc.d directories.
Whatever you do, the server must be run by the PostgreSQL user account and not by root or any other user. Therefore you probably should form your commands using su postgres -c '...'. For example:

```
su postgres -c 'pg_ctl start -D /usr/local/pgsql/data -l serverlog'
```

Here are a few more operating-system-specific suggestions. (In each case be sure to use the proper installation directory and user name where we show generic values.)

- For FreeBSD, look at the file contrib/start-scripts/freebsd in the PostgreSQL source distribution.
- On OpenBSD, add the following lines to the file /etc/rc.local:

```
if [ -x /usr/local/pgsql/bin/pg_ctl -a -x /usr/local/pgsql/bin/ postgres ]; then
  su -l postgres -c '/usr/local/pgsql/bin/pg_ctl start -s -l /var/postgresql/log -D /usr/local/pgsql/data'
  echo -n ' postgresql'
fi
```
- On Linux systems either add

```
/usr/local/pgsql/bin/pg_ctl start -l logfile -D /usr/local/pgsql/data
```
to /etc/rc.d/rc.local or /etc/rc.local or look at the file contrib/start-scripts/linux in the PostgreSQL source distribution.

When using systemd, you can use the following service unit file (e.g., at /etc/systemd/system/postgresql.service):

```
[Unit]
Description=PostgreSQL database server
Documentation=man:postgres(1)

[Service]
Type=notify
User=postgres
ExecStart=/usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data
ExecReload=/bin/kill -HUP $MAINPID
KillMode=mixed
KillSignal=SIGINT
TimeoutSec=0

[Install]
WantedBy=multi-user.target
```

Using Type=notify requires that the server binary was built with configure --with-systemd.

Consider carefully the timeout setting. systemd has a default timeout of 90 seconds as of this writing and will kill a process that does not notify readiness within that time. But a PostgreSQL server that might have to perform crash recovery at startup could take much longer to become ready. The suggested value of 0 disables the timeout logic.

- On NetBSD, use either the FreeBSD or Linux start scripts, depending on preference.
• On Solaris, create a file called `/etc/init.d/postgresql` that contains the following line:

```
su - postgres -c "\"/usr/local/pgsql/bin/pg_ctl start -l logfile -D /usr/local/pgsql/data\""
```

Then, create a symbolic link to it in `/etc/rc3.d` as `S99postgresql`.

While the server is running, its PID is stored in the file `postmaster.pid` in the data directory. This is used to prevent multiple server instances from running in the same data directory and can also be used for shutting down the server.

### 18.3.1. Server Start-up Failures

There are several common reasons the server might fail to start. Check the server's log file, or start it by hand (without redirecting standard output or standard error) and see what error messages appear. Below we explain some of the most common error messages in more detail.

**LOG:** could not bind IPv4 address "127.0.0.1": Address already in use

**HINT:** Is another postmaster already running on port 5432? If not, wait a few seconds and retry.

**FATAL:** could not create any TCP/IP sockets

This usually means just what it suggests: you tried to start another server on the same port where one is already running. However, if the kernel error message is not Address already in use or some variant of that, there might be a different problem. For example, trying to start a server on a reserved port number might draw something like:

```
$ postgres -p 666
```

**LOG:** could not bind IPv4 address "127.0.0.1": Permission denied

**HINT:** Is another postmaster already running on port 666? If not, wait a few seconds and retry.

**FATAL:** could not create any TCP/IP sockets

A message like:

**FATAL:** could not create shared memory segment: Invalid argument

**DETAIL:** Failed system call was shmget(key=5440001, size=4011376640, 03600).

probably means your kernel's limit on the size of shared memory is smaller than the work area PostgreSQL is trying to create (4011376640 bytes in this example). Or it could mean that you do not have System-V-style shared memory support configured into your kernel at all. As a temporary workaround, you can try starting the server with a smaller-than-normal number of buffers (`shared_buffers`). You will eventually want to reconfigure your kernel to increase the allowed shared memory size. You might also see this message when trying to start multiple servers on the same machine, if their total space requested exceeds the kernel limit.

An error like:

**FATAL:** could not create semaphores: No space left on device

**DETAIL:** Failed system call was semget(5440126, 17, 03600).

does not mean you've run out of disk space. It means your kernel's limit on the number of System V semaphores is smaller than the number PostgreSQL wants to create. As above, you might be able...
to work around the problem by starting the server with a reduced number of allowed connections (max_connections), but you'll eventually want to increase the kernel limit.

If you get an “illegal system call” error, it is likely that shared memory or semaphores are not supported in your kernel at all. In that case your only option is to reconfigure the kernel to enable these features.

Details about configuring System V IPC facilities are given in Section 18.4.1.

18.3.2. Client Connection Problems

Although the error conditions possible on the client side are quite varied and application-dependent, a few of them might be directly related to how the server was started. Conditions other than those shown below should be documented with the respective client application.

```
psql: could not connect to server: Connection refused
      Is the server running on host "server.joe.com" and accepting
      TCP/IP connections on port 5432?
```

This is the generic “I couldn't find a server to talk to” failure. It looks like the above when TCP/IP communication is attempted. A common mistake is to forget to configure the server to allow TCP/IP connections.

Alternatively, you'll get this when attempting Unix-domain socket communication to a local server:

```
psql: could not connect to server: No such file or directory
      Is the server running locally and accepting
      connections on Unix domain socket " /tmp/.s.PGSQL.5432"?
```

The last line is useful in verifying that the client is trying to connect to the right place. If there is in fact no server running there, the kernel error message will typically be either Connection refused or No such file or directory, as illustrated. (It is important to realize that Connection refused in this context does not mean that the server got your connection request and rejected it. That case will produce a different message, as shown in Section 20.15.) Other error messages such as Connection timed out might indicate more fundamental problems, like lack of network connectivity.

18.4. Managing Kernel Resources

PostgreSQL can sometimes exhaust various operating system resource limits, especially when multiple copies of the server are running on the same system, or in very large installations. This section explains the kernel resources used by PostgreSQL and the steps you can take to resolve problems related to kernel resource consumption.

18.4.1. Shared Memory and Semaphores

PostgreSQL requires the operating system to provide inter-process communication (IPC) features, specifically shared memory and semaphores. Unix-derived systems typically provide “System V” IPC, “POSIX” IPC, or both. Windows has its own implementation of these features and is not discussed here.

The complete lack of these facilities is usually manifested by an “illegal system call” error upon server start. In that case there is no alternative but to reconfigure your kernel. PostgreSQL won't work without them. This situation is rare, however, among modern operating systems.

Upon starting the server, PostgreSQL normally allocates a very small amount of System V shared memory, as well as a much larger amount of POSIX (mmap) shared memory. In addition a significant
number of semaphores, which can be either System V or POSIX style, are created at server startup. Currently, POSIX semaphores are used on Linux and FreeBSD systems while other platforms use System V semaphores.

Note

Prior to PostgreSQL 9.3, only System V shared memory was used, so the amount of System V shared memory required to start the server was much larger. If you are running an older version of the server, please consult the documentation for your server version.

System V IPC features are typically constrained by system-wide allocation limits. When PostgreSQL exceeds one of these limits, the server will refuse to start and should leave an instructive error message describing the problem and what to do about it. (See also Section 18.3.1.) The relevant kernel parameters are named consistently across different systems; Table 18.1 gives an overview. The methods to set them, however, vary. Suggestions for some platforms are given below.

Table 18.1. System V IPC Parameters

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Values needed to run one PostgreSQL instance</th>
</tr>
</thead>
<tbody>
<tr>
<td>SHMMAX</td>
<td>Maximum size of shared memory segment (bytes)</td>
<td>at least 1kB, but the default is usually much higher</td>
</tr>
<tr>
<td>SHMMIN</td>
<td>Minimum size of shared memory segment (bytes)</td>
<td>1</td>
</tr>
<tr>
<td>SHMALL</td>
<td>Total amount of shared memory available (bytes or pages)</td>
<td>same as SHMMAX if bytes, or ceil(SHMMAX/PAGE_SIZE) if pages, plus room for other applications</td>
</tr>
<tr>
<td>SHMSEG</td>
<td>Maximum number of shared memory segments per process</td>
<td>only 1 segment is needed, but the default is much higher</td>
</tr>
<tr>
<td>SHMMNI</td>
<td>Maximum number of shared memory segments system-wide</td>
<td>like SHMSEG plus room for other applications</td>
</tr>
<tr>
<td>SEMMNI</td>
<td>Maximum number of semaphore identifiers (i.e., sets)</td>
<td>at least ceil((max_connections + autovacuum_max_workers + max_worker_processes + 5) / 16) plus room for other applications</td>
</tr>
<tr>
<td>SEMMNS</td>
<td>Maximum number of semaphores system-wide</td>
<td>ceil((max_connections + autovacuum_max_workers + max_worker_processes + 5) / 16) * 17 plus room for other applications</td>
</tr>
<tr>
<td>SEMMSL</td>
<td>Maximum number of semaphores per set</td>
<td>at least 17</td>
</tr>
<tr>
<td>SEMMAP</td>
<td>Number of entries in semaphore map</td>
<td>see text</td>
</tr>
<tr>
<td>SEMMX</td>
<td>Maximum value of semaphore</td>
<td>at least 1000 (The default is often 32767; do not change unless necessary)</td>
</tr>
</tbody>
</table>
PostgreSQL requires a few bytes of System V shared memory (typically 48 bytes, on 64-bit platforms) for each copy of the server. On most modern operating systems, this amount can easily be allocated. However, if you are running many copies of the server, or if other applications are also using System V shared memory, it may be necessary to increase $\text{SHMALL}$, which is the total amount of System V shared memory system-wide. Note that $\text{SHMALL}$ is measured in pages rather than bytes on many systems.

Less likely to cause problems is the minimum size for shared memory segments ($\text{SHMIN}$), which should be at most approximately 32 bytes for PostgreSQL (it is usually just 1). The maximum number of segments system-wide ($\text{SHMMNI}$) or per-process ($\text{SHMSEG}$) are unlikely to cause a problem unless your system has them set to zero.

When using System V semaphores, PostgreSQL uses one semaphore per allowed connection (max_connections), allowed autovacuum worker process (autovacuum_max_workers) and allowed background process (max_worker_processes), in sets of 16. Each such set will also contain a 17th semaphore which contains a "magic number", to detect collision with semaphore sets used by other applications. The maximum number of semaphores in the system is set by $\text{SEMMNS}$, which consequently must be at least as high as max_connections + autovacuum_max_workers + max_worker_processes, plus one extra for each 16 allowed connections plus workers (see the formula in Table 18.1). The parameter $\text{SEMMNI}$ determines the limit on the number of semaphore sets that can exist on the system at one time. Hence this parameter must be at least $\text{ceil}((\text{max_connections} + \text{autovacuum_max_workers} + \text{max_worker_processes} + 5) / 16)$.

Lowering the number of allowed connections is a temporary workaround for failures, which are usually confusingly worded “No space left on device”, from the function semget.

In some cases it might also be necessary to increase $\text{SEMMAP}$ to be at least on the order of $\text{SEMMNS}$. This parameter defines the size of the semaphore resource map, in which each contiguous block of available semaphores needs an entry. When a semaphore set is freed it is either added to an existing entry that is adjacent to the freed block or it is registered under a new map entry. If the map is full, the freed semaphores get lost (until reboot). Fragmentation of the semaphore space could over time lead to fewer available semaphores than there should be.

Various other settings related to "semaphore undo", such as $\text{SEMMNU}$ and $\text{SEMUME}$, do not affect PostgreSQL.

When using POSIX semaphores, the number of semaphores needed is the same as for System V, that is one semaphore per allowed connection (max_connections), allowed autovacuum worker process (autovacuum_max_workers) and allowed background process (max_worker_processes). On the platforms where this option is preferred, there is no specific kernel limit on the number of POSIX semaphores.

AIX

At least as of version 5.1, it should not be necessary to do any special configuration for such parameters as $\text{SHMMAX}$, as it appears this is configured to allow all memory to be used as shared memory. That is the sort of configuration commonly used for other databases such as DB2.

It might, however, be necessary to modify the global ulimit information in /etc/security/limits, as the default hard limits for file sizes ($\text{fsize}$) and numbers of files ($\text{nofiles}$) might be too low.

FreeBSD

The default settings can be changed using the $\text{sysctl}$ or loader interfaces. The following parameters can be set using $\text{sysctl}$:

```bash
# sysctl kern.ipc.shmall=32768
# sysctl kern.ipc.shmmax=134217728
```

To make these settings persist over reboots, modify /etc/sysctl.conf.
These semaphore-related settings are read-only as far as `sysctl` is concerned, but can be set in `/boot/loader.conf`:

```plaintext
kern.ipc.semmni=256
kern.ipc.semmns=512
kern.ipc.semmnu=256
```

After modifying these values a reboot is required for the new settings to take effect. (Note: FreeBSD does not use SEMMAP. Older versions would accept but ignore a setting for `kern.ipc.semmap`; newer versions reject it altogether.)

You might also want to configure your kernel to lock shared memory into RAM and prevent it from being paged out to swap. This can be accomplished using the `sysctl` setting `kern.ipc.shm_use_phys`.

If running in FreeBSD jails by enabling `sysctl.security.jail.sysvipc_allowed`, postmasters running in different jails should be run by different operating system users. This improves security because it prevents non-root users from interfering with shared memory or semaphores in different jails, and it allows the PostgreSQL IPC cleanup code to function properly. (In FreeBSD 6.0 and later the IPC cleanup code does not properly detect processes in other jails, preventing the running of postmasters on the same port in different jails.)

FreeBSD versions before 4.0 work like OpenBSD (see below).

**NetBSD**

In NetBSD 5.0 and later, IPC parameters can be adjusted using `sysctl`, for example:

```
$ sysctl -w kern.ipc.shmmax=16777216
```

To have these settings persist over reboots, modify `/etc/sysctl.conf`.

You might also want to configure your kernel to lock shared memory into RAM and prevent it from being paged out to swap. This can be accomplished using the `sysctl` setting `kern.ipc.shm_use_phys`.

NetBSD versions before 5.0 work like OpenBSD (see below), except that parameters should be set with the keyword `options` not `option`.

**OpenBSD**

The options `SYSVSHM` and `SYSVSEM` need to be enabled when the kernel is compiled. (They are by default.) The maximum size of shared memory is determined by the option `SHMMAXPGS` (in pages). The following shows an example of how to set the various parameters:

```plaintext
option SYSVSHM
option SHMMAXPGS=4096
option SHMSEG=256

option SYSVSEM
option SEMMNI=256
option SEMMNS=512
option SEMMNU=256
```

**HP-UX**

The default settings tend to suffice for normal installations. On HP-UX 10, the factory default for `SEMMNS` is 128, which might be too low for larger database sites.
IPC parameters can be set in the System Administration Manager (SAM) under Kernel Configuration → Configurable Parameters. Choose Create A New Kernel when you're done.

**Linux**

The default maximum segment size is 32 MB, and the default maximum total size is 2097152 pages. A page is almost always 4096 bytes except in unusual kernel configurations with “huge pages” (use `getconf PAGE_SIZE` to verify).

The shared memory size settings can be changed via the `sysctl` interface. For example, to allow 16 GB:

```
$ sysctl -w kernel.shmmax=17179869184
$ sysctl -w kernel.shmall=4194304
```

In addition these settings can be preserved between reboots in the file `/etc/sysctl.conf`. Doing that is highly recommended.

Ancient distributions might not have the `sysctl` program, but equivalent changes can be made by manipulating the `/proc` file system:

```
$ echo 17179869184 >/proc/sys/kernel/shmmax
$ echo 4194304 >/proc/sys/kernel/shmall
```

The remaining defaults are quite generously sized, and usually do not require changes.

**macOS**

The recommended method for configuring shared memory in macOS is to create a file named `/etc/sysctl.conf`, containing variable assignments such as:

```
kern.sysv.shmmax=4194304
kern.sysv.shmmin=1
kern.sysv.shmmni=32
kern.sysv.shmseg=8
kern.sysv.shmall=1024
```

Note that in some macOS versions, *all five* shared-memory parameters must be set in `/etc/sysctl.conf`, else the values will be ignored.

Beware that recent releases of macOS ignore attempts to set `SHMMAX` to a value that isn’t an exact multiple of 4096.

`SHMALL` is measured in 4 kB pages on this platform.

In older macOS versions, you will need to reboot to have changes in the shared memory parameters take effect. As of 10.5 it is possible to change all but `SHMMNI` on the fly, using `sysctl`. But it’s still best to set up your preferred values via `/etc/sysctl.conf`, so that the values will be kept across reboots.

The file `/etc/sysctl.conf` is only honored in macOS 10.3.9 and later. If you are running a previous 10.3.x release, you must edit the file `/etc/rc` and change the values in the following commands:

```
sysctl -w kern.sysv.shmmax
sysctl -w kern.sysv.shmmin
sysctl -w kern.sysv.shmmni
sysctl -w kern.sysv.shmseg
```
sysctl -w kern.sysv.shmall

Note that /etc/rc is usually overwritten by macOS system updates, so you should expect to have to redo these edits after each update.

In macOS 10.2 and earlier, instead edit these commands in the file /System/Library/StartupItems/SystemTuning/SystemTuning.

Solaris 2.6 to 2.9 (Solaris 6 to Solaris 9)

The relevant settings can be changed in /etc/system, for example:

```
set shmsys:shminfo_shmmax=0x2000000
set shmsys:shminfo_shmmin=1
set shmsys:shminfo_shmmni=256
set shmsys:shminfo_shmseg=256

set semsys:seminfo_semmap=256
set semsys:seminfo_semmni=512
set semsys:seminfo_semmns=512
set semsys:seminfo_semmsl=32
```

You need to reboot for the changes to take effect. See also http://sunsite.uakom.sk/sunworld-online/swol-09-1997/swol-09-09-insidesolaris.html for information on shared memory under older versions of Solaris.

Solaris 2.10 (Solaris 10) and later

OpenSolaris

In Solaris 10 and later, and OpenSolaris, the default shared memory and semaphore settings are good enough for most PostgreSQL applications. Solaris now defaults to a SHMMAX of one-quarter of system RAM. To further adjust this setting, use a project setting associated with the postgres user. For example, run the following as root:

```
projadd -c "PostgreSQL DB User" -K "project.max-shm-memory=(privileged,8GB,deny)" -U postgres -G postgres user.postgres
```

This command adds the user.postgres project and sets the shared memory maximum for the postgres user to 8GB, and takes effect the next time that user logs in, or when you restart PostgreSQL (not reload). The above assumes that PostgreSQL is run by the postgres user in the postgres group. No server reboot is required.

Other recommended kernel setting changes for database servers which will have a large number of connections are:

```
project.max-shm-ids=(priv,32768,deny)
project.max-sem-ids=(priv,4096,deny)
project.max-msg-ids=(priv,4096,deny)
```

Additionally, if you are running PostgreSQL inside a zone, you may need to raise the zone resource usage limits as well. See "Chapter2: Projects and Tasks" in the System Administrator's Guide for more information on projects and prctl.

18.4.2. systemd RemoveIPC

If systemd is in use, some care must be taken that IPC resources (shared memory and semaphores) are not prematurely removed by the operating system. This is especially of concern when installing
PostgreSQL from source. Users of distribution packages of PostgreSQL are less likely to be affected, as the postgres user is then normally created as a system user.

The setting RemoveIPC in logind.conf controls whether IPC objects are removed when a user fully logs out. System users are exempt. This setting defaults to on in stock systemd, but some operating system distributions default it to off.

A typical observed effect when this setting is on is that the semaphore objects used by a PostgreSQL server are removed at apparently random times, leading to the server crashing with log messages like

```
LOG: semctl(1234567890, 0, IPC_RMID, ...) failed: Invalid argument
```

Different types of IPC objects (shared memory vs. semaphores, System V vs. POSIX) are treated slightly differently by systemd, so one might observe that some IPC resources are not removed in the same way as others. But it is not advisable to rely on these subtle differences.

A “user logging out” might happen as part of a maintenance job or manually when an administrator logs in as the postgres user or something similar, so it is hard to prevent in general.

What is a “system user” is determined at systemd compile time from the SYS_UID_MAX setting in /etc/login.defs.

Packaging and deployment scripts should be careful to create the postgres user as a system user by using useradd -r, adduser --system, or equivalent.

Alternatively, if the user account was created incorrectly or cannot be changed, it is recommended to set

```
RemoveIPC=no
```

in /etc/systemd/logind.conf or another appropriate configuration file.

### Caution

At least one of these two things has to be ensured, or the PostgreSQL server will be very unreliable.

#### 18.4.3. Resource Limits

Unix-like operating systems enforce various kinds of resource limits that might interfere with the operation of your PostgreSQL server. Of particular importance are limits on the number of processes per user, the number of open files per process, and the amount of memory available to each process.

Each of these have a “hard” and a “soft” limit. The soft limit is what actually counts but it can be changed by the user up to the hard limit. The hard limit can only be changed by the root user. The system call setrlimit is responsible for setting these parameters. The shell’s built-in command ulimit (Bourne shells) or limit (csh) is used to control the resource limits from the command line.

On BSD-derived systems the file /etc/login.conf controls the various resource limits set during login. See the operating system documentation for details. The relevant parameters are maxproc, openfiles, and datasize. For example:

```
default:\n...
:datasize-cur=256M:\n:maxproc-cur=256:\n:openfiles-cur=256:\n```
Kernels can also have system-wide limits on some resources.

- On Linux `/proc/sys/fs/file-max` determines the maximum number of open files that the kernel will support. It can be changed by writing a different number into the file or by adding an assignment in `/etc/sysctl.conf`. The maximum limit of files per process is fixed at the time the kernel is compiled; see `/usr/src/linux/documentation/proc.txt` for more information.

The PostgreSQL server uses one process per connection so you should provide for at least as many processes as allowed connections, in addition to what you need for the rest of your system. This is usually not a problem but if you run several servers on one machine things might get tight.

The factory default limit on open files is often set to “socially friendly” values that allow many users to coexist on a machine without using an inappropriate fraction of the system resources. If you run many servers on a machine this is perhaps what you want, but on dedicated servers you might want to raise this limit.

On the other side of the coin, some systems allow individual processes to open large numbers of files; if more than a few processes do so then the system-wide limit can easily be exceeded. If you find this happening, and you do not want to alter the system-wide limit, you can set PostgreSQL's `max_files_per_process` configuration parameter to limit the consumption of open files.

### 18.4.4. Linux Memory Overcommit

In Linux 2.4 and later, the default virtual memory behavior is not optimal for PostgreSQL. Because of the way that the kernel implements memory overcommit, the kernel might terminate the PostgreSQL postmaster (the master server process) if the memory demands of either PostgreSQL or another process cause the system to run out of virtual memory.

If this happens, you will see a kernel message that looks like this (consult your system documentation and configuration on where to look for such a message):

```
Out of Memory: Killed process 12345 (postgres).
```

This indicates that the `postgres` process has been terminated due to memory pressure. Although existing database connections will continue to function normally, no new connections will be accepted. To recover, PostgreSQL will need to be restarted.

One way to avoid this problem is to run PostgreSQL on a machine where you can be sure that other processes will not run the machine out of memory. If memory is tight, increasing the swap space of the operating system can help avoid the problem, because the out-of-memory (OOM) killer is invoked only when physical memory and swap space are exhausted.

If PostgreSQL itself is the cause of the system running out of memory, you can avoid the problem by changing your configuration. In some cases, it may help to lower memory-related configuration parameters, particularly `shared_buffers` and `work_mem`. In other cases, the problem may be caused by allowing too many connections to the database server itself. In many cases, it may be better to reduce `max_connections` and instead make use of external connection-pooling software.

On Linux 2.6 and later, it is possible to modify the kernel's behavior so that it will not “overcommit” memory. Although this setting will not prevent the OOM killer\(^1\) from being invoked altogether, it will lower the chances significantly and will therefore lead to more robust system behavior. This is done by selecting strict overcommit mode via `sysctl`:

---

\(^1\) https://lwn.net/Articles/104179/
sysctl -w vm.overcommit_memory=2

or placing an equivalent entry in /etc/sysctl.conf. You might also wish to modify the related setting vm.overcommit_ratio. For details see the kernel documentation file https://www.kernel.org/doc/Documentation/vm/overcommit-accounting.

Another approach, which can be used with or without altering vm.overcommit_memory, is to set the process-specific OOM score adjustment value for the postmaster process to -1000, thereby guaranteeing it will not be targeted by the OOM killer. The simplest way to do this is to execute

echo -1000 > /proc/self/oom_score_adj

in the postmaster's startup script just before invoking the postmaster. Note that this action must be done as root, or it will have no effect; so a root-owned startup script is the easiest place to do it. If you do this, you should also set these environment variables in the startup script before invoking the postmaster:

```bash
export PG_OOM_ADJUST_FILE=/proc/self/oom_score_adj
export PG_OOM_ADJUST_VALUE=0
```

These settings will cause postmaster child processes to run with the normal OOM score adjustment of zero, so that the OOM killer can still target them at need. You could use some other value for PG_OOM_ADJUST_VALUE if you want the child processes to run with some other OOM score adjustment. (PG_OOM_ADJUST_VALUE can also be omitted, in which case it defaults to zero.) If you do not set PG_OOM_ADJUST_FILE, the child processes will run with the same OOM score adjustment as the postmaster, which is unwise since the whole point is to ensure that the postmaster has a preferential setting.

Older Linux kernels do not offer /proc/self/oom_score_adj, but may have a previous version of the same functionality called /proc/self/oom_adj. This works the same except the disable value is -17 not -1000.

---

**Note**

Some vendors' Linux 2.4 kernels are reported to have early versions of the 2.6 overcommit sysctl parameter. However, setting vm.overcommit_memory to 2 on a 2.4 kernel that does not have the relevant code will make things worse, not better. It is recommended that you inspect the actual kernel source code (see the function vm_e-nough_memory in the file mm/mmap.c) to verify what is supported in your kernel before you try this in a 2.4 installation. The presence of the overcommit-accounting documentation file should not be taken as evidence that the feature is there. If in any doubt, consult a kernel expert or your kernel vendor.

---

### 18.4.5. Linux Huge Pages

Using huge pages reduces overhead when using large contiguous chunks of memory, as PostgreSQL does, particularly when using large values of shared_buffers. To use this feature in PostgreSQL you need a kernel with CONFIG_HUGETLBFS=y and CONFIG_HUGETLB_PAGE=y. You will also have to adjust the kernel setting vm.nr_hugepages. To estimate the number of huge pages needed, start PostgreSQL without huge pages enabled and check the postmaster's anonymous shared memory segment size, as well as the system's huge page size, using the /proc file system. This might look like:

```bash
$ head -1 $PGDATA/postmaster.pid
4170
$ pmap 4170 | awk '/rw-s/ && /zero/ {print $2}'
```
$ grep ^Hugepagesize /proc/meminfo
Hugepagesize: 2048 kB

6490428 / 2048 gives approximately 3169.154, so in this example we need at least 3170 huge pages, which we can set with:

$ sysctl -w vm.nr_hugepages=3170

A larger setting would be appropriate if other programs on the machine also need huge pages. Don't forget to add this setting to /etc/sysctl.conf so that it will be reapplied after reboots.

Sometimes the kernel is not able to allocate the desired number of huge pages immediately, so it might be necessary to repeat the command or to reboot. (Immediately after a reboot, most of the machine's memory should be available to convert into huge pages.) To verify the huge page allocation situation, use:

$ grep Huge /proc/meminfo

It may also be necessary to give the database server's operating system user permission to use huge pages by setting vm.hugetlb_shm_group via sysctl, and/or give permission to lock memory with ulimit -l.

The default behavior for huge pages in PostgreSQL is to use them when possible and to fall back to normal pages when failing. To enforce the use of huge pages, you can set huge_pages to on in postgresql.conf. Note that with this setting PostgreSQL will fail to start if not enough huge pages are available.


### 18.5. Shutting Down the Server

There are several ways to shut down the database server. You control the type of shutdown by sending different signals to the master postgres process.

**SIGTERM**

This is the Smart Shutdown mode. After receiving SIGTERM, the server disallows new connections, but lets existing sessions end their work normally. It shuts down only after all of the sessions terminate. If the server is in online backup mode, it additionally waits until online backup mode is no longer active. While backup mode is active, new connections will still be allowed, but only to superusers (this exception allows a superuser to connect to terminate online backup mode). If the server is in recovery when a smart shutdown is requested, recovery and streaming replication will be stopped only after all regular sessions have terminated.

**SIGINT**

This is the Fast Shutdown mode. The server disallows new connections and sends all existing server processes SIGTERM, which will cause them to abort their current transactions and exit promptly. It then waits for all server processes to exit and finally shuts down. If the server is in online backup mode, backup mode will be terminated, rendering the backup useless.

**SIGQUIT**

This is the Immediate Shutdown mode. The server will send SIGQUIT to all child processes and wait for them to terminate. If any do not terminate within 5 seconds, they will be sent SIGKILL. The master server process exits as soon as all child processes have exited, without doing normal
Server Setup and Operation

database shutdown processing. This will lead to recovery (by replaying the WAL log) upon next start-up. This is recommended only in emergencies.

The pg_ctl program provides a convenient interface for sending these signals to shut down the server. Alternatively, you can send the signal directly using kill on non-Windows systems. The PID of the postgres process can be found using the ps program, or from the file postmaster.pid in the data directory. For example, to do a fast shutdown:

```bash
$ kill -INT `head -1 /usr/local/pgsql/data/postmaster.pid`
```

Important

It is best not to use SIGKILL to shut down the server. Doing so will prevent the server from releasing shared memory and semaphores, which might then have to be done manually before a new server can be started. Furthermore, SIGKILL kills the postgres process without letting it relay the signal to its subprocesses, so it will be necessary to kill the individual subprocesses by hand as well.

To terminate an individual session while allowing other sessions to continue, use pg_terminate_backend() (see Table 9.78) or send a SIGTERM signal to the child process associated with the session.

18.6. Upgrading a PostgreSQL Cluster

This section discusses how to upgrade your database data from one PostgreSQL release to a newer one.

Current PostgreSQL version numbers consist of a major and a minor version number. For example, in the version number 10.1, the 10 is the major version number and the 1 is the minor version number, meaning this would be the first minor release of the major release 10. For releases before PostgreSQL version 10.0, version numbers consist of three numbers, for example, 9.5.3. In those cases, the major version consists of the first two digit groups of the version number, e.g., 9.5, and the minor version is the third number, e.g., 3, meaning this would be the third minor release of the major release 9.5.

Minor releases never change the internal storage format and are always compatible with earlier and later minor releases of the same major version number. For example, version 10.1 is compatible with version 10.0 and version 10.6. Similarly, for example, 9.5.3 is compatible with 9.5.0, 9.5.1, and 9.5.6. To update between compatible versions, you simply replace the executables while the server is down and restart the server. The data directory remains unchanged — minor upgrades are that simple.

For major releases of PostgreSQL, the internal data storage format is subject to change, thus complicating upgrades. The traditional method for moving data to a new major version is to dump and reload the database, though this can be slow. A faster method is pg_upgrade. Replication methods are also available, as discussed below.

New major versions also typically introduce some user-visible incompatibilities, so application programming changes might be required. All user-visible changes are listed in the release notes (Appendix E); pay particular attention to the section labeled "Migration". If you are upgrading across several major versions, be sure to read the release notes for each intervening version.

Cautious users will want to test their client applications on the new version before switching over fully; therefore, it's often a good idea to set up concurrent installations of old and new versions. When testing a PostgreSQL major upgrade, consider the following categories of possible changes:

Administration

The capabilities available for administrators to monitor and control the server often change and improve in each major release.
SQL

Typically this includes new SQL command capabilities and not changes in behavior, unless specifically mentioned in the release notes.

Library API

Typically libraries like libpq only add new functionality, again unless mentioned in the release notes.

System Catalogs

System catalog changes usually only affect database management tools.

Server C-language API

This involves changes in the backend function API, which is written in the C programming language. Such changes affect code that references backend functions deep inside the server.

18.6.1. Upgrading Data via pg_dumpall

One upgrade method is to dump data from one major version of PostgreSQL and reload it in another — to do this, you must use a logical backup tool like pg_dumpall; file system level backup methods will not work. (There are checks in place that prevent you from using a data directory with an incompatible version of PostgreSQL, so no great harm can be done by trying to start the wrong server version on a data directory.)

It is recommended that you use the pg_dump and pg_dumpall programs from the newer version of PostgreSQL, to take advantage of enhancements that might have been made in these programs. Current releases of the dump programs can read data from any server version back to 7.0.

These instructions assume that your existing installation is under the /usr/local/pgsql directory, and that the data area is in /usr/local/pgsql/data. Substitute your paths appropriately.

1. If making a backup, make sure that your database is not being updated. This does not affect the integrity of the backup, but the changed data would of course not be included. If necessary, edit the permissions in the file /usr/local/pgsql/data/pg_hba.conf (or equivalent) to disallow access from everyone except you. See Chapter 20 for additional information on access control.

To back up your database installation, type:

`pg_dumpall > outputfile`

To make the backup, you can use the pg_dumpall command from the version you are currently running; see Section 25.1.2 for more details. For best results, however, try to use the pg_dumpall command from PostgreSQL 11.1, since this version contains bug fixes and improvements over older versions. While this advice might seem idiosyncratic since you haven't installed the new version yet, it is advisable to follow it if you plan to install the new version in parallel with the old version. In that case you can complete the installation normally and transfer the data later. This will also decrease the downtime.

2. Shut down the old server:

`pg_ctl stop`

On systems that have PostgreSQL started at boot time, there is probably a start-up file that will accomplish the same thing. For example, on a Red Hat Linux system one might find that this works:
/etc/rc.d/init.d/postgresql stop

See Chapter 18 for details about starting and stopping the server.

3. If restoring from backup, rename or delete the old installation directory if it is not version-specific. It is a good idea to rename the directory, rather than delete it, in case you have trouble and need to revert to it. Keep in mind the directory might consume significant disk space. To rename the directory, use a command like this:

```bash
mv /usr/local/pgsql /usr/local/pgsql.old
```

(Be sure to move the directory as a single unit so relative paths remain unchanged.)

4. Install the new version of PostgreSQL as outlined in Section 16.4.

5. Create a new database cluster if needed. Remember that you must execute these commands while logged in to the special database user account (which you already have if you are upgrading).

```bash
/usr/local/pgsql/bin/initdb -D /usr/local/pgsql/data
```

6. Restore your previous `pg_hba.conf` and any `postgresql.conf` modifications.

7. Start the database server, again using the special database user account:

```bash
/usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data
```

8. Finally, restore your data from backup with:

```bash
/usr/local/pgsql/bin/psql -d postgres -f outputfile
```

using the `new` `psql`.

The least downtime can be achieved by installing the new server in a different directory and running both the old and the new servers in parallel, on different ports. Then you can use something like:

```bash
pg_dumpall -p 5432 | psql -d postgres -p 5433
```

to transfer your data.

### 18.6.2. Upgrading Data via `pg_upgrade`

The `pg_upgrade` module allows an installation to be migrated in-place from one major PostgreSQL version to another. Upgrades can be performed in minutes, particularly with `--link` mode. It requires steps similar to `pg_dumpall` above, e.g. starting/Stopping the server, running `initdb`. The `pg_upgrade` documentation outlines the necessary steps.

### 18.6.3. Upgrading Data via Replication

It is also possible to use logical replication methods to create a standby server with the updated version of PostgreSQL. This is possible because logical replication supports replication between different major versions of PostgreSQL. The standby can be on the same computer or a different computer. Once it has synced up with the master server (running the older version of PostgreSQL), you can switch masters and make the standby the master and shut down the older database instance. Such a switch-over results in only several seconds of downtime for an upgrade.
This method of upgrading can be performed using the built-in logical replication facilities as well as using external logical replication systems such as pglogical, Slony, Londiste, and Bucardo.

### 18.7. Preventing Server Spoofing

While the server is running, it is not possible for a malicious user to take the place of the normal database server. However, when the server is down, it is possible for a local user to spoof the normal server by starting their own server. The spoof server could read passwords and queries sent by clients, but could not return any data because the `PGDATA` directory would still be secure because of directory permissions. Spoofing is possible because any user can start a database server; a client cannot identify an invalid server unless it is specially configured.

One way to prevent spoofing of local connections is to use a Unix domain socket directory (`unix_socket_directories`) that has write permission only for a trusted local user. This prevents a malicious user from creating their own socket file in that directory. If you are concerned that some applications might still reference `/tmp` for the socket file and hence be vulnerable to spoofing, during operating system startup create a symbolic link `/tmp/.s.PGSQL.5432` that points to the relocated socket file. You also might need to modify your `/tmp` cleanup script to prevent removal of the symbolic link.

Another option for local connections is for clients to use `requirepeer` to specify the required owner of the server process connected to the socket.

To prevent spoofing on TCP connections, the best solution is to use SSL certificates and make sure that clients check the server's certificate. To do that, the server must be configured to accept only `hostssl` connections (Section 20.1) and have SSL key and certificate files (Section 18.9). The TCP client must connect using `sslmode=verify-ca` or `verify-full` and have the appropriate root certificate file installed (Section 34.18.1).

### 18.8. Encryption Options

PostgreSQL offers encryption at several levels, and provides flexibility in protecting data from disclosure due to database server theft, unscrupulous administrators, and insecure networks. Encryption might also be required to secure sensitive data such as medical records or financial transactions.

#### Password Encryption

Database user passwords are stored as hashes (determined by the setting `password_encryption`), so the administrator cannot determine the actual password assigned to the user. If SCRAM or MD5 encryption is used for client authentication, the unencrypted password is never even temporarily present on the server because the client encrypts it before being sent across the network. SCRAM is preferred, because it is an Internet standard and is more secure than the PostgreSQL-specific MD5 authentication protocol.

#### Encryption For Specific Columns

The `pgcrypto` module allows certain fields to be stored encrypted. This is useful if only some of the data is sensitive. The client supplies the decryption key and the data is decrypted on the server and then sent to the client.

The decrypted data and the decryption key are present on the server for a brief time while it is being decrypted and communicated between the client and server. This presents a brief moment where the data and keys can be intercepted by someone with complete access to the database server, such as the system administrator.

#### Data Partition Encryption

Storage encryption can be performed at the file system level or the block level. Linux file system encryption options include eCryptfs and EncFS, while FreeBSD uses PEFS. Block level or full
disk encryption options include dm-crypt + LUKS on Linux and GEOM modules geli and gbde on FreeBSD. Many other operating systems support this functionality, including Windows.

This mechanism prevents unencrypted data from being read from the drives if the drives or the entire computer is stolen. This does not protect against attacks while the file system is mounted, because when mounted, the operating system provides an unencrypted view of the data. However, to mount the file system, you need some way for the encryption key to be passed to the operating system, and sometimes the key is stored somewhere on the host that mounts the disk.

Encrypting Data Across A Network

SSL connections encrypt all data sent across the network: the password, the queries, and the data returned. The pg_hba.conf file allows administrators to specify which hosts can use non-encrypted connections (host) and which require SSL-encrypted connections (hostssl). Also, clients can specify that they connect to servers only via SSL. Stunnel or SSH can also be used to encrypt transmissions.

SSL Host Authentication

It is possible for both the client and server to provide SSL certificates to each other. It takes some extra configuration on each side, but this provides stronger verification of identity than the mere use of passwords. It prevents a computer from pretending to be the server just long enough to read the password sent by the client. It also helps prevent “man in the middle” attacks where a computer between the client and server pretends to be the server and reads and passes all data between the client and server.

Client-Side Encryption

If the system administrator for the server's machine cannot be trusted, it is necessary for the client to encrypt the data: this way, unencrypted data never appears on the database server. Data is encrypted on the client before being sent to the server, and database results have to be decrypted on the client before being used.

18.9. Secure TCP/IP Connections with SSL

PostgreSQL has native support for using SSL connections to encrypt client/server communications for increased security. This requires that OpenSSL is installed on both client and server systems and that support in PostgreSQL is enabled at build time (see Chapter 16).

18.9.1. Basic Setup

With SSL support compiled in, the PostgreSQL server can be started with SSL enabled by setting the parameter ssl to on in postgresql.conf. The server will listen for both normal and SSL connections on the same TCP port, and will negotiate with any connecting client on whether to use SSL. By default, this is at the client's option; see Section 20.1 about how to set up the server to require use of SSL for some or all connections.

To start in SSL mode, files containing the server certificate and private key must exist. By default, these files are expected to be named server.crt and server.key, respectively, in the server's data directory, but other names and locations can be specified using the configuration parameters ssl_cert_file and ssl_key_file.

On Unix systems, the permissions on server.key must disallow any access to world or group; achieve this by the command chmod 0600 server.key. Alternatively, the file can be owned by root and have group read access (that is, 0640 permissions). That setup is intended for installations where certificate and key files are managed by the operating system. The user under which the PostgreSQL server runs should then be made a member of the group that has access to those certificate and key files.
If the data directory allows group read access then certificate files may need to be located outside of the data directory in order to conform to the security requirements outlined above. Generally, group access is enabled to allow an unprivileged user to backup the database, and in that case the backup software will not be able to read the certificate files and will likely error.

If the private key is protected with a passphrase, the server will prompt for the passphrase and will not start until it has been entered. Using a passphrase also disables the ability to change the server's SSL configuration without a server restart. Furthermore, passphrase-protected private keys cannot be used at all on Windows.

The first certificate in `server.crt` must be the server's certificate because it must match the server's private key. The certificates of “intermediate” certificate authorities can also be appended to the file. Doing this avoids the necessity of storing intermediate certificates on clients, assuming the root and intermediate certificates were created with `v3_ca` extensions. This allows easier expiration of intermediate certificates.

It is not necessary to add the root certificate to `server.crt`. Instead, clients must have the root certificate of the server's certificate chain.

### 18.9.2. OpenSSL Configuration

PostgreSQL reads the system-wide OpenSSL configuration file. By default, this file is named `openssl.cnf` and is located in the directory reported by `openssl version -d`. This default can be overridden by setting environment variable `OPENSSL_CONF` to the name of the desired configuration file.

OpenSSL supports a wide range of ciphers and authentication algorithms, of varying strength. While a list of ciphers can be specified in the OpenSSL configuration file, you can specify ciphers specifically for use by the database server by modifying `ssl_ciphers` in `postgresql.conf`.

#### Note

It is possible to have authentication without encryption overhead by using `NULL-SHA` or `NULL-MD5` ciphers. However, a man-in-the-middle could read and pass communications between client and server. Also, encryption overhead is minimal compared to the overhead of authentication. For these reasons NULL ciphers are not recommended.

### 18.9.3. Using Client Certificates

To require the client to supply a trusted certificate, place certificates of the root certificate authorities (CAs) you trust in a file in the data directory, set the parameter `ssl_ca_file` in `postgresql.conf` to the new file name, and add the authentication option `clientcert=1` to the appropriate `hostssl` line(s) in `pg_hba.conf`. A certificate will then be requested from the client during SSL connection startup. (See Section 34.18 for a description of how to set up certificates on the client.) The server will verify that the client's certificate is signed by one of the trusted certificate authorities.

Intermediate certificates that chain up to existing root certificates can also appear in the `ssl_ca_file` file if you wish to avoid storing them on clients (assuming the root and intermediate certificates were created with `v3_ca` extensions). Certificate Revocation List (CRL) entries are also checked if the parameter `ssl_crl_file` is set. (See http://h41379.www4.hpe.com/doc/83final/ba554_90007/ch04s02.html for diagrams showing SSL certificate usage.)

The `clientcert` authentication option is available for all authentication methods, but only in `pg_hba.conf` lines specified as `hostssl`. When `clientcert` is not specified or is set to 0, the server will still verify any presented client certificates against its CA file, if one is configured — but it will not insist that a client certificate be presented.
If you are setting up client certificates, you may wish to use the `cert` authentication method, so that the certificates control user authentication as well as providing connection security. See Section 20.12 for details. (It is not necessary to specify `clientcert=1` explicitly when using the `cert` authentication method.)

### 18.9.4. SSL Server File Usage

Table 18.2 summarizes the files that are relevant to the SSL setup on the server. (The shown file names are default names. The locally configured names could be different.)

<table>
<thead>
<tr>
<th>File</th>
<th>Contents</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>ssl_cert_file</td>
<td>server certificate</td>
<td>sent to client to indicate server's identity</td>
</tr>
<tr>
<td>(PGDA-TA/server.crt)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ssl_key_file</td>
<td>server private key</td>
<td>proves server certificate was sent by the owner; does not indicate certificate owner is trustworthy</td>
</tr>
<tr>
<td>(PGDA-TA/server.key)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ssl_ca_file</td>
<td>trusted certificate authorities</td>
<td>checks that client certificate is signed by a trusted certificate authority</td>
</tr>
<tr>
<td>ssl_crl_file</td>
<td>certificates revoked by certificate authorities</td>
<td>client certificate must not be on this list</td>
</tr>
</tbody>
</table>

The server reads these files at server start and whenever the server configuration is reloaded. On Windows systems, they are also re-read whenever a new backend process is spawned for a new client connection.

If an error in these files is detected at server start, the server will refuse to start. But if an error is detected during a configuration reload, the files are ignored and the old SSL configuration continues to be used. On Windows systems, if an error in these files is detected at backend start, that backend will be unable to establish an SSL connection. In all these cases, the error condition is reported in the server log.

### 18.9.5. Creating Certificates

To create a simple self-signed certificate for the server, valid for 365 days, use the following OpenSSL command, replacing `dbhost.yourdomain.com` with the server's host name:

```
openssl req -new -x509 -days 365 -nodes -text -out server.crt -keyout server.key -subj "/CN=dbhost.yourdomain.com"
```

Then do:

```
chmod og-rwx server.key
```

because the server will reject the file if its permissions are more liberal than this. For more details on how to create your server private key and certificate, refer to the OpenSSL documentation.

While a self-signed certificate can be used for testing, a certificate signed by a certificate authority (CA) (usually an enterprise-wide root CA) should be used in production.

To create a server certificate whose identity can be validated by clients, first create a certificate signing request (CSR) and a public/private key file:
openssl req -new -nodes -text -out root.csr \
  -keyout root.key -subj "/CN=root.yourdomain.com"
chmod og-rwx root.key

Then, sign the request with the key to create a root certificate authority (using the default OpenSSL configuration file location on Linux):

openssl x509 -req -in root.csr -text -days 3650 \
  -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
  -signkey root.key -out root.crt

Finally, create a server certificate signed by the new root certificate authority:

openssl req -new -nodes -text -out server.csr \
  -keyout server.key -subj "/CN=dbhost.yourdomain.com"
chmod og-rwx server.key
openssl x509 -req -in server.csr -text -days 365 \
  -CA root.crt -CAkey root.key -CAcreateserial \
  -out server.crt

server.crt and server.key should be stored on the server, and root.crt should be stored on the client so the client can verify that the server's leaf certificate was signed by its trusted root certificate. root.key should be stored offline for use in creating future certificates.

It is also possible to create a chain of trust that includes intermediate certificates:

# root
openssl req -new -nodes -text -out root.csr \
  -keyout root.key -subj "/CN=root.yourdomain.com"
chmod og-rwx root.key
openssl x509 -req -in root.csr -text -days 3650 \
  -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
  -signkey root.key -out root.crt

# intermediate
openssl req -new -nodes -text -out intermediate.csr \
  -keyout intermediate.key -subj "/CN=intermediate.yourdomain.com"
chmod og-rwx intermediate.key
openssl x509 -req -in intermediate.csr -text -days 1825 \
  -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
  -CA root.crt -CAkey root.key -CAcreateserial \
  -out intermediate.crt

# leaf
openssl req -new -nodes -text -out server.csr \
  -keyout server.key -subj "/CN=dbhost.yourdomain.com"
chmod og-rwx server.key
openssl x509 -req -in server.csr -text -days 365 \
  -CA intermediate.crt -CAkey intermediate.key -CAcreateserial \
  -out server.crt

server.crt and intermediate.crt should be concatenated into a certificate file bundle and stored on the server. server.key should also be stored on the server. root.crt should be stored on the client so the client can verify that the server's leaf certificate was signed by a chain of certificates linked to its trusted root certificate. root.key and intermediate.key should be stored offline for use in creating future certificates.
18.10. Secure TCP/IP Connections with SSH Tunnels

It is possible to use SSH to encrypt the network connection between clients and a PostgreSQL server. Done properly, this provides an adequately secure network connection, even for non-SSL-capable clients.

First make sure that an SSH server is running properly on the same machine as the PostgreSQL server and that you can log in using `ssh` as some user. Then you can establish a secure tunnel with a command like this from the client machine:

```
ssh -L 63333:localhost:5432 joe@foo.com
```

The first number in the `-L` argument, 63333, is the port number of your end of the tunnel; it can be any unused port. (IANA reserves ports 49152 through 65535 for private use.) The second number, 5432, is the remote end of the tunnel: the port number your server is using. The name or IP address between the port numbers is the host with the database server you are going to connect to, as seen from the host you are logging in to, which is `foo.com` in this example. In order to connect to the database server using this tunnel, you connect to port 63333 on the local machine:

```
psql -h localhost -p 63333 postgres
```

To the database server it will then look as though you are really user `joe` on host `foo.com` connecting to `localhost` in that context, and it will use whatever authentication procedure was configured for connections from this user and host. Note that the server will not think the connection is SSL-encrypted, since in fact it is not encrypted between the SSH server and the PostgreSQL server. This should not pose any extra security risk as long as they are on the same machine.

In order for the tunnel setup to succeed you must be allowed to connect via `ssh` as `joe@foo.com`, just as if you had attempted to use `ssh` to create a terminal session.

You could also have set up the port forwarding as

```
ssh -L 63333:foo.com:5432 joe@foo.com
```

but then the database server will see the connection as coming in on its `foo.com` interface, which is not opened by the default setting `listen_addresses = 'localhost'`. This is usually not what you want.

If you have to “hop” to the database server via some login host, one possible setup could look like this:

```
ssh -L 63333:db.foo.com:5432 joe@shell.foo.com
```

Note that this way the connection from `shell.foo.com` to `db.foo.com` will not be encrypted by the SSH tunnel. SSH offers quite a few configuration possibilities when the network is restricted in various ways. Please refer to the SSH documentation for details.

**Tip**

Several other applications exist that can provide secure tunnels using a procedure similar in concept to the one just described.
18.11. Registering Event Log on Windows

To register a Windows event log library with the operating system, issue this command:

```
regsvr32 pgsql_library_directory/pgevent.dll
```

This creates registry entries used by the event viewer, under the default event source named PostgreSQL.

To specify a different event source name (see event_source), use the /n and /i options:

```
regsvr32 /n /i:event_source_name pgsql_library_directory/pgevent.dll
```

To unregister the event log library from the operating system, issue this command:

```
regsvr32 /u [/i:event_source_name] pgsql_library_directory/pgevent.dll
```

Note

To enable event logging in the database server, modify log_destination to include eventlog in postgresql.conf.
Chapter 19. Server Configuration

There are many configuration parameters that affect the behavior of the database system. In the first section of this chapter we describe how to interact with configuration parameters. The subsequent sections discuss each parameter in detail.

19.1. Setting Parameters

19.1.1. Parameter Names and Values

All parameter names are case-insensitive. Every parameter takes a value of one of five types: boolean, string, integer, floating point, or enumerated (enum). The type determines the syntax for setting the parameter:

- **Boolean**: Values can be written as `on`, `off`, `true`, `false`, `yes`, `no`, `1`, `0` (all case-insensitive) or any unambiguous prefix of one of these.

- **String**: In general, enclose the value in single quotes, doubling any single quotes within the value. Quotes can usually be omitted if the value is a simple number or identifier, however.

- **Numeric (integer and floating point)**: A decimal point is permitted only for floating-point parameters. Do not use thousands separators. Quotes are not required.

- **Numeric with Unit**: Some numeric parameters have an implicit unit, because they describe quantities of memory or time. The unit might be bytes, kilobytes, blocks (typically eight kilobytes), milliseconds, seconds, or minutes. An unadorned numeric value for one of these settings will use the setting's default unit, which can be learned from `pg_settings.unit`. For convenience, settings can be given with a unit specified explicitly, for example `'120 ms'` for a time value, and they will be converted to whatever the parameter's actual unit is. Note that the value must be written as a string (with quotes) to use this feature. The unit name is case-sensitive, and there can be whitespace between the numeric value and the unit.

  - Valid memory units are B (bytes), kB (kilobytes), MB (megabytes), GB (gigabytes), and TB (terabytes). The multiplier for memory units is 1024, not 1000.

  - Valid time units are ms (milliseconds), s (seconds), min (minutes), h (hours), and d (days).

- **Enumerated**: Enumerated-type parameters are written in the same way as string parameters, but are restricted to have one of a limited set of values. The values allowable for such a parameter can be found from `pg_settings.enumvals`. Enum parameter values are case-insensitive.

19.1.2. Parameter Interaction via the Configuration File

The most fundamental way to set these parameters is to edit the file `postgresql.conf`, which is normally kept in the data directory. A default copy is installed when the database cluster directory is initialized. An example of what this file might look like is:

```plaintext
# This is a comment
log_connections = yes
log_destination = 'syslog'
search_path = '"$user", public'
shared_buffers = 128MB
```

One parameter is specified per line. The equal sign between name and value is optional. Whitespace is insignificant (except within a quoted parameter value) and blank lines are ignored. Hash marks (`#`) designate the remainder of the line as a comment. Parameter values that are not simple identifiers or
numbers must be single-quoted. To embed a single quote in a parameter value, write either two quotes (preferred) or backslash-quote.

Parameters set in this way provide default values for the cluster. The settings seen by active sessions will be these values unless they are overridden. The following sections describe ways in which the administrator or user can override these defaults.

The configuration file is reread whenever the main server process receives a SIGHUP signal; this signal is most easily sent by running `pg_ctl reload` from the command line or by calling the SQL function `pg_reload_conf()`. The main server process also propagates this signal to all currently running server processes, so that existing sessions also adopt the new values (this will happen after they complete any currently-executing client command). Alternatively, you can send the signal to a single server process directly. Some parameters can only be set at server start; any changes to their entries in the configuration file will be ignored until the server is restarted. Invalid parameter settings in the configuration file are likewise ignored (but logged) during SIGHUP processing.

In addition to `postgresql.conf`, a PostgreSQL data directory contains a file `postgresql.auto.conf`, which has the same format as `postgresql.conf` but should never be edited manually. This file holds settings provided through the ALTER SYSTEM command. This file is automatically read whenever `postgresql.conf` is, and its settings take effect in the same way. Settings in `postgresql.auto.conf` override those in `postgresql.conf`.

The system view `pg_file_settings` can be helpful for pre-testing changes to the configuration file, or for diagnosing problems if a SIGHUP signal did not have the desired effects.

### 19.1.3. Parameter Interaction via SQL

PostgreSQL provides three SQL commands to establish configuration defaults. The already-mentioned ALTER SYSTEM command provides a SQL-accessible means of changing global defaults; it is functionally equivalent to editing `postgresql.conf`. In addition, there are two commands that allow setting of defaults on a per-database or per-role basis:

- The ALTER DATABASE command allows global settings to be overridden on a per-database basis.
- The ALTER ROLE command allows both global and per-database settings to be overridden with user-specific values.

Values set with ALTER DATABASE and ALTER ROLE are applied only when starting a fresh database session. They override values obtained from the configuration files or server command line, and constitute defaults for the rest of the session. Note that some settings cannot be changed after server start, and so cannot be set with these commands (or the ones listed below).

Once a client is connected to the database, PostgreSQL provides two additional SQL commands (and equivalent functions) to interact with session-local configuration settings:

- The SHOW command allows inspection of the current value of all parameters. The corresponding function is `current_setting(setting_name text)`.
- The SET command allows modification of the current value of those parameters that can be set locally to a session; it has no effect on other sessions. The corresponding function is `set_config(setting_name, new_value, is_local)`.

In addition, the system view `pg_settings` can be used to view and change session-local values:

- Querying this view is similar to using SHOW ALL but provides more detail. It is also more flexible, since it's possible to specify filter conditions or join against other relations.
- Using UPDATE on this view, specifically updating the setting column, is the equivalent of issuing SET commands. For example, the equivalent of
SET configuration_parameter TO DEFAULT;

is:

UPDATE pg_settings SET setting = reset_val WHERE name = 'configuration_parameter';

19.1.4. Parameter Interaction via the Shell

In addition to setting global defaults or attaching overrides at the database or role level, you can pass settings to PostgreSQL via shell facilities. Both the server and libpq client library accept parameter values via the shell.

• During server startup, parameter settings can be passed to the `postgres` command via the `-c` command-line parameter. For example,

```
postgres -c log_connections=yes -c log_destination=syslog
```

Settings provided in this way override those set via `postgresql.conf` or `ALTER SYSTEM`, so they cannot be changed globally without restarting the server.

• When starting a client session via libpq, parameter settings can be specified using the `PGOPTIONS` environment variable. Settings established in this way constitute defaults for the life of the session, but do not affect other sessions. For historical reasons, the format of `PGOPTIONS` is similar to that used when launching the `postgres` command; specifically, the `-c` flag must be specified. For example,

```
env PGOPTIONS="-c geqo=off -c statement_timeout=5min" psql
```

Other clients and libraries might provide their own mechanisms, via the shell or otherwise, that allow the user to alter session settings without direct use of SQL commands.

19.1.5. Managing Configuration File Contents

PostgreSQL provides several features for breaking down complex `postgresql.conf` files into sub-files. These features are especially useful when managing multiple servers with related, but not identical, configurations.

In addition to individual parameter settings, the `postgresql.conf` file can contain `include` directives, which specify another file to read and process as if it were inserted into the configuration file at this point. This feature allows a configuration file to be divided into physically separate parts. Include directives simply look like:

```
include 'filename'
```

If the file name is not an absolute path, it is taken as relative to the directory containing the referencing configuration file. Inclusions can be nested.

There is also an `include_if_exists` directive, which acts the same as the `include` directive, except when the referenced file does not exist or cannot be read. A regular `include` will consider this an error condition, but `include_if_exists` merely logs a message and continues processing the referencing configuration file.

The `postgresql.conf` file can also contain `include_dir` directives, which specify an entire directory of configuration files to include. These look like
Server Configuration

include_dir 'directory'

Non-absolute directory names are taken as relative to the directory containing the referencing configuration file. Within the specified directory, only non-directory files whose names end with the suffix .conf will be included. File names that start with the . character are also ignored, to prevent mistakes since such files are hidden on some platforms. Multiple files within an include directory are processed in file name order (according to C locale rules, i.e. numbers before letters, and uppercase letters before lowercase ones).

Include files or directories can be used to logically separate portions of the database configuration, rather than having a single large postgresql.conf file. Consider a company that has two database servers, each with a different amount of memory. There are likely elements of the configuration both will share, for things such as logging. But memory-related parameters on the server will vary between the two. And there might be server specific customizations, too. One way to manage this situation is to break the custom configuration changes for your site into three files. You could add this to the end of your postgresql.conf file to include them:

include 'shared.conf'
include 'memory.conf'
include 'server.conf'

All systems would have the same shared.conf. Each server with a particular amount of memory could share the same memory.conf; you might have one for all servers with 8GB of RAM, another for those having 16GB. And finally server.conf could have truly server-specific configuration information in it.

Another possibility is to create a configuration file directory and put this information into files there. For example, a conf.d directory could be referenced at the end of postgresql.conf:

include_dir 'conf.d'

Then you could name the files in the conf.d directory like this:

00shared.conf
01memory.conf
02server.conf

This naming convention establishes a clear order in which these files will be loaded. This is important because only the last setting encountered for a particular parameter while the server is reading configuration files will be used. In this example, something set in conf.d/02server.conf would override a value set in conf.d/01memory.conf.

You might instead use this approach to naming the files descriptively:

00shared.conf
01memory-8GB.conf
02server-foo.conf

This sort of arrangement gives a unique name for each configuration file variation. This can help eliminate ambiguity when several servers have their configurations all stored in one place, such as in a version control repository. (Storing database configuration files under version control is another good practice to consider.)

19.2. File Locations

In addition to the postgresql.conf file already mentioned, PostgreSQL uses two other manually-edited configuration files, which control client authentication (their use is discussed in Chapter 20).
By default, all three configuration files are stored in the database cluster's data directory. The parameters described in this section allow the configuration files to be placed elsewhere. (Doing so can ease administration. In particular it is often easier to ensure that the configuration files are properly backed-up when they are kept separate.)

**data_directory** (string)

Specifies the directory to use for data storage. This parameter can only be set at server start.

**config_file** (string)

Specifies the main server configuration file (customarily called `postgresql.conf`). This parameter can only be set on the `postgres` command line.

**hba_file** (string)

Specifies the configuration file for host-based authentication (customarily called `pg_hba.conf`). This parameter can only be set at server start.

**ident_file** (string)

Specifies the configuration file for user name mapping (customarily called `pg_ident.conf`). This parameter can only be set at server start. See also Section 20.2.

**external_pid_file** (string)

Specifies the name of an additional process-ID (PID) file that the server should create for use by server administration programs. This parameter can only be set at server start.

In a default installation, none of the above parameters are set explicitly. Instead, the data directory is specified by the `-D` command-line option or the `PGDATA` environment variable, and the configuration files are all found within the data directory.

If you wish to keep the configuration files elsewhere than the data directory, the `postgres -D` command-line option or `PGDATA` environment variable must point to the directory containing the configuration files, and the `data_directory` parameter must be set in `postgresql.conf` (or on the command line) to show where the data directory is actually located. Notice that `data_directory` overrides `-D` and `PGDATA` for the location of the data directory, but not for the location of the configuration files.

If you wish, you can specify the configuration file names and locations individually using the parameters `config_file`, `hba_file` and/or `ident_file`. `config_file` can only be specified on the `postgres` command line, but the others can be set within the main configuration file. If all three parameters plus `data_directory` are explicitly set, then it is not necessary to specify `-D` or `PGDATA`.

When setting any of these parameters, a relative path will be interpreted with respect to the directory in which `postgres` is started.

### 19.3. Connections and Authentication

#### 19.3.1. Connection Settings

**listen_addresses** (string)

Specifies the TCP/IP address(es) on which the server is to listen for connections from client applications. The value takes the form of a comma-separated list of host names and/or numeric IP addresses. The special entry `*` corresponds to all available IP interfaces. The entry `0.0.0.0` allows listening for all IPv4 addresses and `::` allows listening for all IPv6 addresses. If the list is empty, the server does not listen on any IP interface at all, in which case only Unix-domain sockets can be used to connect to it. The default value is `localhost`, which allows only local TCP/IP
“loopback” connections to be made. While client authentication (Chapter 20) allows fine-grained control over who can access the server, listen_addresses controls which interfaces accept connection attempts, which can help prevent repeated malicious connection requests on insecure network interfaces. This parameter can only be set at server start.

port (integer)

The TCP port the server listens on; 5432 by default. Note that the same port number is used for all IP addresses the server listens on. This parameter can only be set at server start.

max_connections (integer)

Determines the maximum number of concurrent connections to the database server. The default is typically 100 connections, but might be less if your kernel settings will not support it (as determined during initdb). This parameter can only be set at server start.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

superuser_reserved_connections (integer)

Determines the number of connection “slots” that are reserved for connections by PostgreSQL superusers. At most max_connections connections can ever be active simultaneously. Whenever the number of active concurrent connections is at least max_connections minus superuser_reserved_connections, new connections will be accepted only for superusers, and no new replication connections will be accepted.

The default value is three connections. The value must be less than max_connections minus max_wal_senders. This parameter can only be set at server start.

unix_socket_directories (string)

Specifies the directory of the Unix-domain socket(s) on which the server is to listen for connections from client applications. Multiple sockets can be created by listing multiple directories separated by commas. Whitespace between entries is ignored; surround a directory name with double quotes if you need to include whitespace or commas in the name. An empty value specifies not listening on any Unix-domain sockets, in which case only TCP/IP sockets can be used to connect to the server. The default value is normally /tmp, but that can be changed at build time. This parameter can only be set at server start.

In addition to the socket file itself, which is named .s.PGSQL.nnnn where nnnn is the server’s port number, an ordinary file named .s.PGSQL.nnnn.lock will be created in each of the unix_socket_directories directories. Neither file should ever be removed manually.

This parameter is irrelevant on Windows, which does not have Unix-domain sockets.

unix_socket_group (string)

Sets the owning group of the Unix-domain socket(s). (The owning user of the sockets is always the user that starts the server.) In combination with the parameter unix_socket_permissions this can be used as an additional access control mechanism for Unix-domain connections. By default this is the empty string, which uses the default group of the server user. This parameter can only be set at server start.

This parameter is irrelevant on Windows, which does not have Unix-domain sockets.

unix_socket_permissions (integer)

Sets the access permissions of the Unix-domain socket(s). Unix-domain sockets use the usual Unix file system permission set. The parameter value is expected to be a numeric mode specified in the format accepted by the chmod and umask system calls. (To use the customary octal format the number must start with a 0 (zero).)

534
The default permissions are 0777, meaning anyone can connect. Reasonable alternatives are 0770 (only user and group, see also unix_socket_group) and 0700 (only user). (Note that for a Unix-domain socket, only write permission matters, so there is no point in setting or revoking read or execute permissions.)

This access control mechanism is independent of the one described in Chapter 20.

This parameter can only be set at server start.

This parameter is irrelevant on systems, notably Solaris as of Solaris 10, that ignore socket permissions entirely. There, one can achieve a similar effect by pointing unix_socket_directories to a directory having search permission limited to the desired audience. This parameter is also irrelevant on Windows, which does not have Unix-domain sockets.

**bonjour (boolean)**

Enables advertising the server's existence via Bonjour. The default is off. This parameter can only be set at server start.

**bonjour_name (string)**

Specifies the Bonjour service name. The computer name is used if this parameter is set to the empty string '' (which is the default). This parameter is ignored if the server was not compiled with Bonjour support. This parameter can only be set at server start.

**tcp_keepalives_idle (integer)**

Specifies the number of seconds of inactivity after which TCP should send a keepalive message to the client. A value of 0 uses the system default. This parameter is supported only on systems that support TCP_KEEPIDLE or an equivalent socket option, and on Windows; on other systems, it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and always reads as zero.

**Note**

On Windows, a value of 0 will set this parameter to 2 hours, since Windows does not provide a way to read the system default value.

**tcp_keepalives_interval (integer)**

Specifies the number of seconds after which a TCP keepalive message that is not acknowledged by the client should be retransmitted. A value of 0 uses the system default. This parameter is supported only on systems that support TCP_KEEPINTVL or an equivalent socket option, and on Windows; on other systems, it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and always reads as zero.

**Note**

On Windows, a value of 0 will set this parameter to 1 second, since Windows does not provide a way to read the system default value.

**tcp_keepalives_count (integer)**

Specifies the number of TCP keepalives that can be lost before the server's connection to the client is considered dead. A value of 0 uses the system default. This parameter is supported only on systems that support TCP_KEEPCNT or an equivalent socket option; on other systems, it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and always reads as zero.
19.3.2. Authentication

**authentication_timeout (integer)**

Maximum time to complete client authentication, in seconds. If a would-be client has not completed the authentication protocol in this much time, the server closes the connection. This prevents hung clients from occupying a connection indefinitely. The default is one minute (1m). This parameter can only be set in the `postgresql.conf` file or on the server command line.

**password_encryption (enum)**

When a password is specified in CREATE ROLE or ALTER ROLE, this parameter determines the algorithm to use to encrypt the password. The default value is `md5`, which stores the password as an MD5 hash (`on` is also accepted, as alias for `md5`). Setting this parameter to `scram-sha-256` will encrypt the password with SCRAM-SHA-256.

Note that older clients might lack support for the SCRAM authentication mechanism, and hence not work with passwords encrypted with SCRAM-SHA-256. See Section 20.5 for more details.

**krb_server_keyfile (string)**

Sets the location of the Kerberos server key file. See Section 20.6 for details. This parameter can only be set in the `postgresql.conf` file or on the server command line.

**krb_caseins_users (boolean)**

Sets whether GSSAPI user names should be treated case-insensitively. The default is `off` (case sensitive). This parameter can only be set in the `postgresql.conf` file or on the server command line.

**db_user_namespace (boolean)**

This parameter enables per-database user names. It is off by default. This parameter can only be set in the `postgresql.conf` file or on the server command line.

If this is on, you should create users as `username@dbname`. When `username` is passed by a connecting client, `@` and the database name are appended to the user name and that database-specific user name is looked up by the server. Note that when you create users with names containing `@` within the SQL environment, you will need to quote the user name.

With this parameter enabled, you can still create ordinary global users. Simply append `@` when specifying the user name in the client, e.g. `joe@`. The `@` will be stripped off before the user name is looked up by the server.

`db_user_namespace` causes the client's and server's user name representation to differ. Authentication checks are always done with the server's user name so authentication methods must be configured for the server's user name, not the client's. Because `md5` uses the user name as salt on both the client and server, `md5` cannot be used with `db_user_namespace`.

**Note**

This feature is intended as a temporary measure until a complete solution is found. At that time, this option will be removed.
19.3.3. SSL

See Section 18.9 for more information about setting up SSL.

**ssl (boolean)**

Enables SSL connections. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default is **off**.

**ssl_ca_file (string)**

Specifies the name of the file containing the SSL server certificate authority (CA). Relative paths are relative to the data directory. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default is empty, meaning no CA file is loaded, and client certificate verification is not performed.

**ssl_cert_file (string)**

Specifies the name of the file containing the SSL server certificate. Relative paths are relative to the data directory. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default is *server.crt*.

**ssl_crl_file (string)**

Specifies the name of the file containing the SSL server certificate revocation list (CRL). Relative paths are relative to the data directory. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default is empty, meaning no CRL file is loaded.

**ssl_key_file (string)**

Specifies the name of the file containing the SSL server private key. Relative paths are relative to the data directory. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default is *server.key*.

**ssl_ciphers (string)**

Specifies a list of SSL cipher suites that are allowed to be used on secure connections. See the ciphers manual page in the OpenSSL package for the syntax of this setting and a list of supported values. This parameter can only be set in the *postgresql.conf* file or on the server command line. The default value is **HIGH:MEDIUM:+3DES:!aNULL**. The default is usually a reasonable choice unless you have specific security requirements.

Explanation of the default value:

**HIGH**

Cipher suites that use ciphers from **HIGH** group (e.g., AES, Camellia, 3DES)

**MEDIUM**

Cipher suites that use ciphers from **MEDIUM** group (e.g., RC4, SEED)

**+3DES**

The OpenSSL default order for **HIGH** is problematic because it orders 3DES higher than AES128. This is wrong because 3DES offers less security than AES128, and it is also much slower. **+3DES** reorders it after all other **HIGH** and **MEDIUM** ciphers.

**!aNULL**

Disables anonymous cipher suites that do no authentication. Such cipher suites are vulnerable to man-in-the-middle attacks and therefore should not be used.
Available cipher suite details will vary across OpenSSL versions. Use the command `openssl ciphers -v 'HIGH:MEDIUM:+3DES:!aNULL'` to see actual details for the currently installed OpenSSL version. Note that this list is filtered at run time based on the server key type.

`ssl_prefer_server_ciphers (boolean)`

Specifies whether to use the server's SSL cipher preferences, rather than the client's. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is `true`.

Older PostgreSQL versions do not have this setting and always use the client's preferences. This setting is mainly for backward compatibility with those versions. Using the server's preferences is usually better because it is more likely that the server is appropriately configured.

`ssl_ecdh_curve (string)`

Specifies the name of the curve to use in ECDH key exchange. It needs to be supported by all clients that connect. It does not need to be the same curve used by the server's Elliptic Curve key. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is `prime256v1`.

OpenSSL names for the most common curves are: `prime256v1` (NIST P-256), `secp384r1` (NIST P-384), `secp521r1` (NIST P-521). The full list of available curves can be shown with the command `openssl ecparam -list_curves`. Not all of them are usable in TLS though.

`ssl_dh_params_file (string)`

Specifies the name of the file containing Diffie-Hellman parameters used for so-called ephemeral DH family of SSL ciphers. The default is empty, in which case compiled-in default DH parameters used. Using custom DH parameters reduces the exposure if an attacker manages to crack the well-known compiled-in DH parameters. You can create your own DH parameters file with the command `openssl dhparam -out dhparams.pem 2048`.

This parameter can only be set in the `postgresql.conf` file or on the server command line.

`ssl_passphrase_command (string)`

Sets an external command to be invoked when a passphrase for decrypting an SSL file such as a private key needs to be obtained. By default, this parameter is empty, which means the built-in prompting mechanism is used.

The command must print the passphrase to the standard output and exit with code 0. In the parameter value, `%p` is replaced by a prompt string. (Write `%` for a literal `%`.) Note that the prompt string will probably contain whitespace, so be sure to quote adequately. A single newline is stripped from the end of the output if present.

The command does not actually have to prompt the user for a passphrase. It can read it from a file, obtain it from a keychain facility, or similar. It is up to the user to make sure the chosen mechanism is adequately secure.

This parameter can only be set in the `postgresql.conf` file or on the server command line.

`ssl_passphrase_command_supports_reload (boolean)`

This parameter determines whether the passphrase command set by `ssl_passphrase_command` will also be called during a configuration reload if a key file needs a passphrase. If this parameter is false (the default), then `ssl_passphrase_command` will be ignored during a reload and the SSL configuration will not be reloaded if a passphrase is needed. That setting is appropriate for a command that requires a TTY for prompting, which might not be available when the server is running. Setting this parameter to true might be appropriate if the passphrase is obtained from a file, for example.
This parameter can only be set in the `postgresql.conf` file or on the server command line.

## 19.4. Resource Consumption

### 19.4.1. Memory

**shared_buffers** (integer)

Sets the amount of memory the database server uses for shared memory buffers. The default is typically 128 megabytes (128MB), but might be less if your kernel settings will not support it (as determined during initdb). This setting must be at least 128 kilobytes. (Non-default values of `BLCKSZ` change the minimum.) However, settings significantly higher than the minimum are usually needed for good performance. This parameter can only be set at server start.

If you have a dedicated database server with 1GB or more of RAM, a reasonable starting value for `shared_buffers` is 25% of the memory in your system. There are some workloads where even larger settings for `shared_buffers` are effective, but because PostgreSQL also relies on the operating system cache, it is unlikely that an allocation of more than 40% of RAM to `shared_buffers` will work better than a smaller amount. Larger settings for `shared_buffers` usually require a corresponding increase in `max_wal_size`, in order to spread out the process of writing large quantities of new or changed data over a longer period of time.

On systems with less than 1GB of RAM, a smaller percentage of RAM is appropriate, so as to leave adequate space for the operating system.

**huge_pages** (enum)

Controls whether huge pages are requested for the main shared memory area. Valid values are `try` (the default), `on`, and `off`. With `huge_pages` set to `try`, the server will try to request huge pages, but fall back to the default if that fails. With `on`, failure to request huge pages will prevent the server from starting up. With `off`, huge pages will not be requested.

At present, this setting is supported only on Linux and Windows. The setting is ignored on other systems when set to `try`.

The use of huge pages results in smaller page tables and less CPU time spent on memory management, increasing performance. For more details about using huge pages on Linux, see Section 18.4.5.

Huge pages are known as large pages on Windows. To use them, you need to assign the user right Lock Pages in Memory to the Windows user account that runs PostgreSQL. You can use Windows Group Policy tool (`gpedit.msc`) to assign the user right Lock Pages in Memory. To start the database server on the command prompt as a standalone process, not as a Windows service, the command prompt must be run as an administrator or User Access Control (UAC) must be disabled. When the UAC is enabled, the normal command prompt revokes the user right Lock Pages in Memory when started.

Note that this setting only affects the main shared memory area. Operating systems such as Linux, FreeBSD, and Illumos can also use huge pages (also known as “super” pages or “large” pages) automatically for normal memory allocation, without an explicit request from PostgreSQL. On Linux, this is called “transparent huge pages” (THP). That feature has been known to cause performance degradation with PostgreSQL for some users on some Linux versions, so its use is currently discouraged (unlike explicit use of `huge_pages`).

**temp_buffers** (integer)

Sets the maximum number of temporary buffers used by each database session. These are session-local buffers used only for access to temporary tables. The default is eight megabytes (8MB).
The setting can be changed within individual sessions, but only before the first use of temporary tables within the session; subsequent attempts to change the value will have no effect on that session.

A session will allocate temporary buffers as needed up to the limit given by `temp_buffers`. The cost of setting a large value in sessions that do not actually need many temporary buffers is only a buffer descriptor, or about 64 bytes, per increment in `temp_buffers`. However if a buffer is actually used an additional 8192 bytes will be consumed for it (or in general, `BLCKSZ` bytes).

`max_prepared_transactions(integer)`

Sets the maximum number of transactions that can be in the "prepared" state simultaneously (see `PREPARE TRANSACTION`). Setting this parameter to zero (which is the default) disables the prepared-transaction feature. This parameter can only be set at server start.

If you are not planning to use prepared transactions, this parameter should be set to zero to prevent accidental creation of prepared transactions. If you are using prepared transactions, you will probably want `max_prepared_transactions` to be at least as large as `max_connections`, so that every session can have a prepared transaction pending.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

`work_mem(integer)`

Specifies the amount of memory to be used by internal sort operations and hash tables before writing to temporary disk files. The value defaults to four megabytes (4MB). Note that for a complex query, several sort or hash operations might be running in parallel; each operation will be allowed to use as much memory as this value specifies before it starts to write data into temporary files. Also, several running sessions could be doing such operations concurrently. Therefore, the total memory used could be many times the value of `work_mem`; it is necessary to keep this fact in mind when choosing the value. Sort operations are used for `ORDER BY`, `DISTINCT`, and merge joins. Hash tables are used in hash joins, hash-based aggregation, and hash-based processing of `IN` subqueries.

`maintenance_work_mem(integer)`

Specifies the maximum amount of memory to be used by maintenance operations, such as `VACUUM`, `CREATE INDEX`, and `ALTER TABLE ADD FOREIGN KEY`. It defaults to 64 megabytes (64MB). Since only one of these operations can be executed at a time by a database session, and an installation normally doesn't have many of them running concurrently, it's safe to set this value significantly larger than `work_mem`. Larger settings might improve performance for vacuuming and for restoring database dumps.

Note that when autovacuum runs, up to `autovacuum_max_workers` times this memory may be allocated, so be careful not to set the default value too high. It may be useful to control for this by separately setting `autovacuum_work_mem`.

`autovacuum_work_mem(integer)`

Specifies the maximum amount of memory to be used by each autovacuum worker process. It defaults to -1, indicating that the value of `maintenance_work_mem` should be used instead. The setting has no effect on the behavior of `VACUUM` when run in other contexts.

`max_stack_depth(integer)`

Specifies the maximum safe depth of the server's execution stack. The ideal setting for this parameter is the actual stack size limit enforced by the kernel (as set by `ulimit -s` or local equivalent), less a safety margin of a megabyte or so. The safety margin is needed because the stack
depth is not checked in every routine in the server, but only in key potentially-recursive routines such as expression evaluation. The default setting is two megabytes (2MB), which is conservatively small and unlikely to risk crashes. However, it might be too small to allow execution of complex functions. Only superusers can change this setting.

Setting max_stack_depth higher than the actual kernel limit will mean that a runaway recursive function can crash an individual backend process. On platforms where PostgreSQL can determine the kernel limit, the server will not allow this variable to be set to an unsafe value. However, not all platforms provide the information, so caution is recommended in selecting a value.

dynamic_shared_memory_type(enum)

Specifies the dynamic shared memory implementation that the server should use. Possible values are posix (for POSIX shared memory allocated using shm_open), sysv (for System V shared memory allocated via shmget), windows (for Windows shared memory), mmap (to simulate shared memory using memory-mapped files stored in the data directory), and none (to disable this feature). Not all values are supported on all platforms; the first supported option is the default for that platform. The use of the mmap option, which is not the default on any platform, is generally discouraged because the operating system may write modified pages back to disk repeatedly, increasing system I/O load; however, it may be useful for debugging, when the pg_dynshmem directory is stored on a RAM disk, or when other shared memory facilities are not available.

19.4.2. Disk

temp_file_limit(integer)

Specifies the maximum amount of disk space that a process can use for temporary files, such as sort and hash temporary files, or the storage file for a held cursor. A transaction attempting to exceed this limit will be canceled. The value is specified in kilobytes, and -1 (the default) means no limit. Only superusers can change this setting.

This setting constrains the total space used at any instant by all temporary files used by a given PostgreSQL process. It should be noted that disk space used for explicit temporary tables, as opposed to temporary files used behind-the-scenes in query execution, does not count against this limit.

19.4.3. Kernel Resource Usage

max_files_per_process(integer)

Sets the maximum number of simultaneously open files allowed to each server subprocess. The default is one thousand files. If the kernel is enforcing a safe per-process limit, you don’t need to worry about this setting. But on some platforms (notably, most BSD systems), the kernel will allow individual processes to open many more files than the system can actually support if many processes all try to open that many files. If you find yourself seeing “Too many open files” failures, try reducing this setting. This parameter can only be set at server start.

19.4.4. Cost-based Vacuum Delay

During the execution of VACUUM and ANALYZE commands, the system maintains an internal counter that keeps track of the estimated cost of the various I/O operations that are performed. When the accumulated cost reaches a limit (specified by vacuum_cost_limit), the process performing the operation will sleep for a short period of time, as specified by vacuum_cost_delay. Then it will reset the counter and continue execution.

The intent of this feature is to allow administrators to reduce the I/O impact of these commands on concurrent database activity. There are many situations where it is not important that maintenance commands like VACUUM and ANALYZE finish quickly; however, it is usually very important that...
these commands do not significantly interfere with the ability of the system to perform other database operations. Cost-based vacuum delay provides a way for administrators to achieve this.

This feature is disabled by default for manually issued VACUUM commands. To enable it, set the vacuum_cost_delay variable to a nonzero value.

vacuum_cost_delay(integer)

The length of time, in milliseconds, that the process will sleep when the cost limit has been exceeded. The default value is zero, which disables the cost-based vacuum delay feature. Positive values enable cost-based vacuuming. Note that on many systems, the effective resolution of sleep delays is 10 milliseconds; setting vacuum_cost_delay to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10.

When using cost-based vacuuming, appropriate values for vacuum_cost_delay are usually quite small, perhaps 10 or 20 milliseconds. Adjusting vacuum's resource consumption is best done by changing the other vacuum cost parameters.

vacuum_cost_page_hit(integer)

The estimated cost for vacuuming a buffer found in the shared buffer cache. It represents the cost to lock the buffer pool, lookup the shared hash table and scan the content of the page. The default value is one.

vacuum_cost_page_miss(integer)

The estimated cost for vacuuming a buffer that has to be read from disk. This represents the effort to lock the buffer pool, lookup the shared hash table, read the desired block in from the disk and scan its content. The default value is 10.

vacuum_cost_page_dirty(integer)

The estimated cost charged when vacuum modifies a block that was previously clean. It represents the extra I/O required to flush the dirty block out to disk again. The default value is 20.

vacuum_cost_limit(integer)

The accumulated cost that will cause the vacuuming process to sleep. The default value is 200.

There are certain operations that hold critical locks and should therefore complete as quickly as possible. Cost-based vacuum delays do not occur during such operations. Therefore it is possible that the cost accumulates far higher than the specified limit. To avoid uselessly long delays in such cases, the actual delay is calculated as vacuum_cost_delay * accumulated_balance / vacuum_cost_limit with a maximum of vacuum_cost_delay * 4.

19.4.5. Background Writer

There is a separate server process called the background writer, whose function is to issue writes of “dirty” (new or modified) shared buffers. It writes shared buffers so server processes handling user queries seldom or never need to wait for a write to occur. However, the background writer does cause a net overall increase in I/O load, because while a repeatedly-dirtied page might otherwise be written only once per checkpoint interval, the background writer might write it several times as it is dirtied in the same interval. The parameters discussed in this subsection can be used to tune the behavior for local needs.
Server Configuration

bgwriter_delay(integer)

Specifies the delay between activity rounds for the background writer. In each round the writer issues writes for some number of dirty buffers (controllable by the following parameters). It then sleeps for `bgwriter_delay` milliseconds, and repeats. When there are no dirty buffers in the buffer pool, though, it goes into a longer sleep regardless of `bgwriter_delay`. The default value is 200 milliseconds (200ms). Note that on many systems, the effective resolution of sleep delays is 10 milliseconds; setting `bgwriter_delay` to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10. This parameter can only be set in the `postgresql.conf` file or on the server command line.

bgwriter_lru_maxpages(integer)

In each round, no more than this many buffers will be written by the background writer. Setting this to zero disables background writing. (Note that checkpoints, which are managed by a separate, dedicated auxiliary process, are unaffected.) The default value is 100 buffers. This parameter can only be set in the `postgresql.conf` file or on the server command line.

bgwriter_lru_multiplier(floating point)

The number of dirty buffers written in each round is based on the number of new buffers that have been needed by server processes during recent rounds. The average recent need is multiplied by `bgwriter_lru_multiplier` to arrive at an estimate of the number of buffers that will be needed during the next round. Dirty buffers are written until there are that many clean, reusable buffers available. (However, no more than `bgwriter_lru_maxpages` buffers will be written per round.) Thus, a setting of 1.0 represents a “just in time” policy of writing exactly the number of buffers predicted to be needed. Larger values provide some cushion against spikes in demand, while smaller values intentionally leave writes to be done by server processes. The default is 2.0. This parameter can only be set in the `postgresql.conf` file or on the server command line.

bgwriter_flush_after(integer)

Whenever more than `bgwriter_flush_after` bytes have been written by the background writer, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel's page cache, reducing the likelihood of stalls when an `fsync` is issued at the end of a checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than `shared_buffers`, but smaller than the OS's page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 512kB on Linux, 0 elsewhere. (If `BLCKSZ` is not 8kB, the default and maximum values scale proportionally to it.) This parameter can only be set in the `postgresql.conf` file or on the server command line.

Smaller values of `bgwriter_lru_maxpages` and `bgwriter_lru_multiplier` reduce the extra I/O load caused by the background writer, but make it more likely that server processes will have to issue writes for themselves, delaying interactive queries.

19.4.6. Asynchronous Behavior

effective_io_concurrency(integer)

Sets the number of concurrent disk I/O operations that PostgreSQL expects can be executed simultaneously. Raising this value will increase the number of I/O operations that any individual PostgreSQL session attempts to initiate in parallel. The allowed range is 1 to 1000, or zero to disable issuance of asynchronous I/O requests. Currently, this setting only affects bitmap heap scans.

For magnetic drives, a good starting point for this setting is the number of separate drives comprising a RAID 0 stripe or RAID 1 mirror being used for the database. (For RAID 5 the parity drive should not be counted.) However, if the database is often busy with multiple queries issued...
in concurrent sessions, lower values may be sufficient to keep the disk array busy. A value higher
than needed to keep the disks busy will only result in extra CPU overhead. SSDs and other
memory-based storage can often process many concurrent requests, so the best value might be
in the hundreds.

Asynchronous I/O depends on an effective `posix_fadvise` function, which some operating
systems lack. If the function is not present then setting this parameter to anything but zero will
result in an error. On some operating systems (e.g., Solaris), the function is present but does not
actually do anything.

The default is 1 on supported systems, otherwise 0. This value can be overridden for tables in a
particular tablespace by setting the tablespace parameter of the same name (see ALTER TABLESPACE).

**max_worker_processes** (integer)

Sets the maximum number of background processes that the system can support. This parameter
can only be set at server start. The default is 8.

When running a standby server, you must set this parameter to the same or higher value than on
the master server. Otherwise, queries will not be allowed in the standby server.

When changing this value, consider also adjusting `max_parallel_workers`, `max_parallel_maintenance_workers`, and `max_parallel_workers_per_gather`.

**max_parallel_workers_per_gather** (integer)

Sets the maximum number of workers that can be started by a single Gather or Gather Merge
node. Parallel workers are taken from the pool of processes established by `max_worker_processes`,
limited by `max_parallel_workers`. Note that the requested number of workers may not actually
be available at run time. If this occurs, the plan will run with fewer workers than expected, which
may be inefficient. The default value is 2. Setting this value to 0 disables parallel query execution.

Note that parallel queries may consume very substantially more resources than non-parallel
queries, because each worker process is a completely separate process which has roughly the
same impact on the system as an additional user session. This should be taken into account when
choosing a value for this setting, as well as when configuring other settings that control resource
utilization, such as `work_mem`. Resource limits such as `work_mem` are applied individually to
each worker, which means the total utilization may be much higher across all processes than it
would normally be for any single process. For example, a parallel query using 4 workers may use
up to 5 times as much CPU time, memory, I/O bandwidth, and so forth as a query which uses
no workers at all.

For more information on parallel query, see Chapter 15.

**max_parallel_maintenance_workers** (integer)

Sets the maximum number of parallel workers that can be started by a single utility command. Currently, the only parallel utility command that supports the use of parallel workers is `CREATE INDEX`, and only when building a B-tree index. Parallel workers are taken from the pool of processes established by `max_worker_processes`, limited by `max_parallel_workers`. Note that the requested number of workers may not actually be available at run time. If this occurs, the utility operation will run with fewer workers than expected. The default value is 2. Setting this value to 0 disables the use of parallel workers by utility commands.

Note that parallel utility commands should not consume substantially more memory than equiva-

tent non-parallel operations. This strategy differs from that of parallel query, where resource lim-
its generally apply per worker process. Parallel utility commands treat the resource limit `main-
tenance_work_mem` as a limit to be applied to the entire utility command, regardless of the
number of parallel worker processes. However, parallel utility commands may still consume sub-
stantially more CPU resources and I/O bandwidth.
Server Configuration

max_parallel_workers(integer)

Sets the maximum number of workers that the system can support for parallel operations. The default value is 8. When increasing or decreasing this value, consider also adjusting max_parallel_maintenance_workers and max_parallel_workers_per_gather. Also, note that a setting for this value which is higher than max_worker_processes will have no effect, since parallel workers are taken from the pool of worker processes established by that setting.

backend_flush_after(integer)

Whenever more than backend_flush_after bytes have been written by a single backend, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel's page cache, reducing the likelihood of stalls when an fsync is issued at the end of a checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than shared_buffers, but smaller than the OS's page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 0, i.e., no forced writeback. (If BLCKSZ is not 8kB, the maximum value scales proportionally to it.)

old_snapshot_threshold(integer)

Sets the minimum time that a snapshot can be used without risk of a snapshot too old error occurring when using the snapshot. This parameter can only be set at server start.

Beyond the threshold, old data may be vacuumed away. This can help prevent bloat in the face of snapshots which remain in use for a long time. To prevent incorrect results due to cleanup of data which would otherwise be visible to the snapshot, an error is generated when the snapshot is older than this threshold and the snapshot is used to read a page which has been modified since the snapshot was built.

A value of -1 disables this feature, and is the default. Useful values for production work probably range from a small number of hours to a few days. The setting will be coerced to a granularity of minutes, and small numbers (such as 0 or 1min) are only allowed because they may sometimes be useful for testing. While a setting as high as 60d is allowed, please note that in many workloads extreme bloat or transaction ID wraparound may occur in much shorter time frames.

When this feature is enabled, freed space at the end of a relation cannot be released to the operating system, since that could remove information needed to detect the snapshot too old condition. All space allocated to a relation remains associated with that relation for reuse only within that relation unless explicitly freed (for example, with VACUUM FULL).

This setting does not attempt to guarantee that an error will be generated under any particular circumstances. In fact, if the correct results can be generated from (for example) a cursor which has materialized a result set, no error will be generated even if the underlying rows in the referenced table have been vacuumed away. Some tables cannot safely be vacuumed early, and so will not be affected by this setting, such as system catalogs. For such tables this setting will neither reduce bloat nor create a possibility of a snapshot too old error on scanning.

19.5. Write Ahead Log

For additional information on tuning these settings, see Section 30.4.

19.5.1. Settings

wal_level(enum)

wal_level determines how much information is written to the WAL. The default value is replica, which writes enough data to support WAL archiving and replication, including run-
ning read-only queries on a standby server. minimal removes all logging except the information required to recover from a crash or immediate shutdown. Finally, logical adds information necessary to support logical decoding. Each level includes the information logged at all lower levels. This parameter can only be set at server start.

In minimal level, WAL-logging of some bulk operations can be safely skipped, which can make those operations much faster (see Section 14.4.7). Operations in which this optimization can be applied include:

CREATE TABLE AS
CREATE INDEX
CLUSTER
COPY into tables that were created or truncated in the same transaction

But minimal WAL does not contain enough information to reconstruct the data from a base backup and the WAL logs, so replica or higher must be used to enable WAL archiving (archive_mode) and streaming replication.

In logical level, the same information is logged as with replica, plus information needed to allow extracting logical change sets from the WAL. Using a level of logical will increase the WAL volume, particularly if many tables are configured for REPLICA IDENTITY FULL and many UPDATE and DELETE statements are executed.

In releases prior to 9.6, this parameter also allowed the values archive and hot_standby. These are still accepted but mapped to replica.

fsync (boolean)

If this parameter is on, the PostgreSQL server will try to make sure that updates are physically written to disk, by issuing fsync() system calls or various equivalent methods (see wal_sync_method). This ensures that the database cluster can recover to a consistent state after an operating system or hardware crash.

While turning off fsync is often a performance benefit, this can result in unrecoverable data corruption in the event of a power failure or system crash. Thus it is only advisable to turn off fsync if you can easily recreate your entire database from external data.

Examples of safe circumstances for turning off fsync include the initial loading of a new database cluster from a backup file, using a database cluster for processing a batch of data after which the database will be thrown away and recreated, or for a read-only database clone which gets recreated frequently and is not used for failover. High quality hardware alone is not a sufficient justification for turning off fsync.

For reliable recovery when changing fsync off to on, it is necessary to force all modified buffers in the kernel to durable storage. This can be done while the cluster is shutdown or while fsync is on by running initdb --sync-only, running sync, unmounting the file system, or rebooting the server.

In many situations, turning off synchronous_commit for noncritical transactions can provide much of the potential performance benefit of turning off fsync, without the attendant risks of data corruption.

fsync can only be set in the postgresql.conf file or on the server command line. If you turn this parameter off, also consider turning off full_page_writes.

synchronous_commit (enum)

Specifies whether transaction commit will wait for WAL records to be written to disk before the command returns a "success" indication to the client. Valid values are on, remote_apply, remote_write, local, and off. The default, and safe, setting is on. When off, there can be a delay between when success is reported to the client and when the transaction is really guar-
anteed to be safe against a server crash. (The maximum delay is three times wal_writer_delay.) Unlike fsync, setting this parameter to off does not create any risk of database inconsistency: an operating system or database crash might result in some recent allegedly-committed transactions being lost, but the database state will be just the same as if those transactions had been aborted cleanly. So, turning synchronous_commit off can be a useful alternative when performance is more important than exact certainty about the durability of a transaction. For more discussion see Section 30.3.

If synchronous_standby_names is non-empty, this parameter also controls whether or not transaction commits will wait for their WAL records to be replicated to the standby server(s). When set to on, commits will wait until replies from the current synchronous standby(s) indicate they have received the commit record of the transaction and flushed it to disk. This ensures the transaction will not be lost unless both the primary and all synchronous standbys suffer corruption of their database storage. When set to remote_apply, commits will wait until replies from the current synchronous standby(s) indicate they have received the commit record of the transaction and applied it, so that it has become visible to queries on the standby(s). When set to remote_write, commits will wait until replies from the current synchronous standby(s) indicate they have received the commit record of the transaction and written it out to their operating system. This setting is sufficient to ensure data preservation even if a standby instance of PostgreSQL were to crash, but not if the standby suffers an operating-system-level crash, since the data has not necessarily reached stable storage on the standby. Finally, the setting local causes commits to wait for local flush to disk, but not for replication. This is not usually desirable when synchronous replication is in use, but is provided for completeness.

If synchronous_standby_names is empty, the settings on, remote_apply, remote_write and local all provide the same synchronization level: transaction commits only wait for local flush to disk.

This parameter can be changed at any time; the behavior for any one transaction is determined by the setting in effect when it commits. It is therefore possible, and useful, to have some transactions commit synchronously and others asynchronously. For example, to make a single multistatement transaction commit asynchronously when the default is the opposite, issue SET LOCAL synchronous_commit TO OFF within the transaction.

wal_sync_method (enum)

Method used for forcing WAL updates out to disk. If fsync is off then this setting is irrelevant, since WAL file updates will not be forced out at all. Possible values are:

• open_datasync (write WAL files with open() option O_DSYNC)
• fdatasync (call fdatasync() at each commit)
• fsync (call fsync() at each commit)
• fsync_writethrough (call fsync() at each commit, forcing write-through of any disk write cache)
• open_sync (write WAL files with open() option O_SYNC)

The open_* options also use O_DIRECT if available. Not all of these choices are available on all platforms. The default is the first method in the above list that is supported by the platform, except that fdatasync is the default on Linux. The default is not necessarily ideal; it might be necessary to change this setting or other aspects of your system configuration in order to create a crash-safe configuration or achieve optimal performance. These aspects are discussed in Section 30.1. This parameter can only be set in the postgresql.conf file or on the server command line.

full_page_writes (boolean)

When this parameter is on, the PostgreSQL server writes the entire content of each disk page to WAL during the first modification of that page after a checkpoint. This is needed because a
page write that is in process during an operating system crash might be only partially completed, leading to an on-disk page that contains a mix of old and new data. The row-level change data normally stored in WAL will not be enough to completely restore such a page during post-crash recovery. Storing the full page image guarantees that the page can be correctly restored, but at the price of increasing the amount of data that must be written to WAL. (Because WAL replay always starts from a checkpoint, it is sufficient to do this during the first change of each page after a checkpoint. Therefore, one way to reduce the cost of full-page writes is to increase the checkpoint interval parameters.)

Turning this parameter off speeds normal operation, but might lead to either unrecoverable data corruption, or silent data corruption, after a system failure. The risks are similar to turning off fsync, though smaller, and it should be turned off only based on the same circumstances recommended for that parameter.

Turning off this parameter does not affect use of WAL archiving for point-in-time recovery (PITR) (see Section 25.3).

This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is on.

**wal_log_hints (boolean)**

When this parameter is on, the PostgreSQL server writes the entire content of each disk page to WAL during the first modification of that page after a checkpoint, even for non-critical modifications of so-called hint bits.

If data checksums are enabled, hint bit updates are always WAL-logged and this setting is ignored. You can use this setting to test how much extra WAL-logging would occur if your database had data checksums enabled.

This parameter can only be set at server start. The default value is off.

**wal_compression (boolean)**

When this parameter is on, the PostgreSQL server compresses a full page image written to WAL when `full_page_writes` is on or during a base backup. A compressed page image will be decompressed during WAL replay. The default value is off. Only superusers can change this setting.

Turning this parameter on can reduce the WAL volume without increasing the risk of unrecoverable data corruption, but at the cost of some extra CPU spent on the compression during WAL logging and on the decompression during WAL replay.

**wal_buffers (integer)**

The amount of shared memory used for WAL data that has not yet been written to disk. The default setting of -1 selects a size equal to 1/32nd (about 3%) of `shared_buffers`, but not less than 64kB nor more than the size of one WAL segment, typically 16MB. This value can be set manually if the automatic choice is too large or too small, but any positive value less than 32kB will be treated as 32kB. This parameter can only be set at server start.

The contents of the WAL buffers are written out to disk at every transaction commit, so extremely large values are unlikely to provide a significant benefit. However, setting this value to at least a few megabytes can improve write performance on a busy server where many clients are committing at once. The auto-tuning selected by the default setting of -1 should give reasonable results in most cases.

**wal_writer_delay (integer)**

Specifies how often the WAL writer flushes WAL. After flushing WAL it sleeps for `wal_writer_delay` milliseconds, unless woken up by an asynchronously committing transaction. If the last flush happened less than `wal_writer_delay` milliseconds ago and less
Server Configuration

than wal_writer_flush_after bytes of WAL have been produced since, then WAL is only written to the operating system, not flushed to disk. The default value is 200 milliseconds (200ms). Note that on many systems, the effective resolution of sleep delays is 10 milliseconds; setting wal_writer_delay to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10. This parameter can only be set in the postgresql.conf file or on the server command line.

wal_writer_flush_after(integer)

Specifies how often the WAL writer flushes WAL. If the last flush happened less than wal_writer_delay milliseconds ago and less than wal_writer_flush_after bytes of WAL have been produced since, then WAL is only written to the operating system, not flushed to disk. If wal_writer_flush_after is set to 0 then WAL data is flushed immediately. The default is 1MB. This parameter can only be set in the postgresql.conf file or on the server command line.

commit_delay(integer)

commit_delay adds a time delay, measured in microseconds, before a WAL flush is initiated. This can improve group commit throughput by allowing a larger number of transactions to commit via a single WAL flush, if system load is high enough that additional transactions become ready to commit within the given interval. However, it also increases latency by up to commit_delay microseconds for each WAL flush. Because the delay is just wasted if no other transactions become ready to commit, a delay is only performed if at least commit_siblings other transactions are active when a flush is about to be initiated. Also, no delays are performed if fsync is disabled. The default commit_delay is zero (no delay). Only superusers can change this setting.

In PostgreSQL releases prior to 9.3, commit_delay behaved differently and was much less effective: it affected only commits, rather than all WAL flushes, and waited for the entire configured delay even if the WAL flush was completed sooner. Beginning in PostgreSQL 9.3, the first process that becomes ready to flush waits for the configured interval, while subsequent processes wait only until the leader completes the flush operation.

commit_siblings(integer)

Minimum number of concurrent open transactions to require before performing the commit_delay delay. A larger value makes it more probable that at least one other transaction will become ready to commit during the delay interval. The default is five transactions.

19.5.2. Checkpoints

checkpoint_timeout(integer)

Maximum time between automatic WAL checkpoints, in seconds. The valid range is between 30 seconds and one day. The default is five minutes (5min). Increasing this parameter can increase the amount of time needed for crash recovery. This parameter can only be set in the postgresql.conf file or on the server command line.

checkpoint_completion_target(floating point)

Specifies the target of checkpoint completion, as a fraction of total time between checkpoints. The default is 0.5. This parameter can only be set in the postgresql.conf file or on the server command line.

checkpoint_flush_after(integer)

Whenever more than checkpoint_flush_after bytes have been written while performing a checkpoint, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel's page cache, reducing the likelihood of stalls when an
fsync is issued at the end of the checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than shared_buffers, but smaller than the OS's page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 256kB on Linux, 0 elsewhere. (If BLCKSZ is not 8kB, the default and maximum values scale proportionally to it.) This parameter can only be set in the `postgresql.conf` file or on the server command line.

**checkpoint_warning (integer)**

Write a message to the server log if checkpoints caused by the filling of WAL segment files happen closer together than this many seconds (which suggests that max_wal_size ought to be raised). The default is 30 seconds (30s). Zero disables the warning. No warnings will be generated if checkpoint_timeout is less than checkpoint_warning. This parameter can only be set in the `postgresql.conf` file or on the server command line.

**max_wal_size (integer)**

Maximum size to let the WAL grow to between automatic WAL checkpoints. This is a soft limit; WAL size can exceed max_wal_size under special circumstances, like under heavy load, a failing `archive_command`, or a high `wal_keep_segments` setting. The default is 1 GB. Increasing this parameter can increase the amount of time needed for crash recovery. This parameter can only be set in the `postgresql.conf` file or on the server command line.

**min_wal_size (integer)**

As long as WAL disk usage stays below this setting, old WAL files are always recycled for future use at a checkpoint, rather than removed. This can be used to ensure that enough WAL space is reserved to handle spikes in WAL usage, for example when running large batch jobs. The default is 80 MB. This parameter can only be set in the `postgresql.conf` file or on the server command line.

### 19.5.3. Archiving

**archive_mode (enum)**

When `archive_mode` is enabled, completed WAL segments are sent to archive storage by setting `archive_command`. In addition to off, to disable, there are two modes: on, and always. During normal operation, there is no difference between the two modes, but when set to always the WAL archiver is enabled also during archive recovery or standby mode. In always mode, all files restored from the archive or streamed with streaming replication will be archived (again). See Section 26.2.9 for details.

`archive_mode` and `archive_command` are separate variables so that `archive_command` can be changed without leaving archiving mode. This parameter can only be set at server start. `archive_mode` cannot be enabled when `wal_level` is set to minimal.

**archive_command (string)**

The local shell command to execute to archive a completed WAL file segment. Any %p in the string is replaced by the path name of the file to archive, and any %f is replaced by only the file name. (The path name is relative to the working directory of the server, i.e., the cluster's data directory.) Use % to embed an actual % character in the command. It is important for the command to return a zero exit status only if it succeeds. For more information see Section 25.3.1.

This parameter can only be set in the `postgresql.conf` file or on the server command line. It is ignored unless `archive_mode` was enabled at server start. If `archive_command` is an empty string (the default) while `archive_mode` is enabled, WAL archiving is temporarily disabled, but the server continues to accumulate WAL segment files in the expectation that a
command will soon be provided. Setting archive_command to a command that does nothing but return true, e.g. /bin/true (REM on Windows), effectively disables archiving, but also breaks the chain of WAL files needed for archive recovery, so it should only be used in unusual circumstances.

archive_timeout (integer)

The archive_command is only invoked for completed WAL segments. Hence, if your server generates little WAL traffic (or has slack periods where it does so), there could be a long delay between the completion of a transaction and its safe recording in archive storage. To limit how old unarchived data can be, you can set archive_timeout to force the server to switch to a new WAL segment file periodically. When this parameter is greater than zero, the server will switch to a new segment file whenever this many seconds have elapsed since the last segment file switch, and there has been any database activity, including a single checkpoint (checkpoints are skipped if there is no database activity). Note that archived files that are closed early due to a forced switch are still the same length as completely full files. Therefore, it is unwise to use a very short archive_timeout — it will bloat your archive storage. archive_timeout settings of a minute or so are usually reasonable. You should consider using streaming replication, instead of archiving, if you want data to be copied off the master server more quickly than that. This parameter can only be set in the postgresql.conf file or on the server command line.

19.6. Replication

These settings control the behavior of the built-in streaming replication feature (see Section 26.2.5). Servers will be either a master or a standby server. Masters can send data, while standbys are always receivers of replicated data. When cascading replication (see Section 26.2.7) is used, standby servers can also be senders, as well as receivers. Parameters are mainly for sending and standby servers, though some parameters have meaning only on the master server. Settings may vary across the cluster without problems if that is required.

19.6.1. Sending Servers

These parameters can be set on any server that is to send replication data to one or more standby servers. The master is always a sending server, so these parameters must always be set on the master. The role and meaning of these parameters does not change after a standby becomes the master.

max_wal_senders (integer)

Specifies the maximum number of concurrent connections from standby servers or streaming base backup clients (i.e., the maximum number of simultaneously running WAL sender processes). The default is 10. The value 0 means replication is disabled. WAL sender processes count towards the total number of connections, so this parameter's value must be less than max_connections minus superuser_reserved_connections. Abrupt streaming client disconnection might leave an orphaned connection slot behind until a timeout is reached, so this parameter should be set slightly higher than the maximum number of expected clients so disconnected clients can immediately reconnect. This parameter can only be set at server start. Also, wal_level must be set to replica or higher to allow connections from standby servers.

max_replication_slots (integer)

Specifies the maximum number of replication slots (see Section 26.2.6) that the server can support. The default is 10. This parameter can only be set at server start. Setting it to a lower value than the number of currently existing replication slots will prevent the server from starting. Also, wal_level must be set to replica or higher to allow replication slots to be used.

wal_keep_segments (integer)

Specifies the minimum number of past log file segments kept in the pg_wal directory, in case a standby server needs to fetch them for streaming replication. Each segment is normally
16 megabytes. If a standby server connected to the sending server falls behind by more than `wal_keep_segments` segments, the sending server might remove a WAL segment still needed by the standby, in which case the replication connection will be terminated. Downstream connections will also eventually fail as a result. (However, the standby server can recover by fetching the segment from archive, if WAL archiving is in use.)

This sets only the minimum number of segments retained in `pg_wal`; the system might need to retain more segments for WAL archival or to recover from a checkpoint. If `wal_keep_segments` is zero (the default), the system doesn't keep any extra segments for standby purposes, so the number of old WAL segments available to standby servers is a function of the location of the previous checkpoint and status of WAL archiving. This parameter can only be set in the `postgresql.conf` file or on the server command line.

### `wal_sender_timeout` (integer)

Terminate replication connections that are inactive longer than the specified number of milliseconds. This is useful for the sending server to detect a standby crash or network outage. A value of zero disables the timeout mechanism. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default value is 60 seconds.

### `track_commit_timestamp` (boolean)

Record commit time of transactions. This parameter can only be set in `postgresql.conf` file or on the server command line. The default value is `off`.

## 19.6.2. Master Server

These parameters can be set on the master/primary server that is to send replication data to one or more standby servers. Note that in addition to these parameters, `wal_level` must be set appropriately on the master server, and optionally WAL archiving can be enabled as well (see Section 19.5.3). The values of these parameters on standby servers are irrelevant, although you may wish to set them there in preparation for the possibility of a standby becoming the master.

### `synchronous_standby_names` (string)

Specifies a list of standby servers that can support synchronous replication, as described in Section 26.2.8. There will be one or more active synchronous standbys; transactions waiting for commit will be allowed to proceed after these standby servers confirm receipt of their data. The synchronous standbys will be those whose names appear in this list, and that are both currently connected and streaming data in real-time (as shown by a state of `streaming` in the `pg_stat_replication` view). Specifying more than one synchronous standby can allow for very high availability and protection against data loss.

The name of a standby server for this purpose is the `application_name` setting of the standby, as set in the standby's connection information. In case of a physical replication standby, this should be set in the `primary_conninfo` setting in `recovery.conf`; the default is `wal_receiver`. For logical replication, this can be set in the connection information of the subscription, and it defaults to the subscription name. For other replication stream consumers, consult their documentation.

This parameter specifies a list of standby servers using either of the following syntaxes:

```
[FIRST] num_sync ( standby_name [, ...] )
ANY num_sync ( standby_name [, ...] )
standby_name [, ...]
```

where `num_sync` is the number of synchronous standbys that transactions need to wait for replies from, and `standby_name` is the name of a standby server. `FIRST` and `ANY` specify the method to choose synchronous standbys from the listed servers.
The keyword FIRST, coupled with num_sync, specifies a priority-based synchronous replication and makes transaction commits wait until their WAL records are replicated to num_sync synchronous standbys chosen based on their priorities. For example, a setting of FIRST 3 (s1, s2, s3, s4) will cause each commit to wait for replies from three higher-priority standbys chosen from standby servers s1, s2, s3 and s4. The standbys whose names appear earlier in the list are given higher priority and will be considered as synchronous. Other standby servers appearing later in this list represent potential synchronous standbys. If any of the current synchronous standbys disconnects for whatever reason, it will be replaced immediately with the next-highest-priority standby. The keyword FIRST is optional.

The keyword ANY, coupled with num_sync, specifies a quorum-based synchronous replication and makes transaction commits wait until their WAL records are replicated to at least num_sync listed standbys. For example, a setting of ANY 3 (s1, s2, s3, s4) will cause each commit to proceed as soon as at least any three standbys of s1, s2, s3 and s4 reply.

FIRST and ANY are case-insensitive. If these keywords are used as the name of a standby server, its standby_name must be double-quoted.

The third syntax was used before PostgreSQL version 9.6 and is still supported. It's the same as the first syntax with FIRST and num_sync equal to 1. For example, FIRST 1 (s1, s2) and s1, s2 have the same meaning: either s1 or s2 is chosen as a synchronous standby.

The special entry * matches any standby name.

There is no mechanism to enforce uniqueness of standby names. In case of duplicates one of the matching standbys will be considered as higher priority, though exactly which one is indeterminate.

**Note**

Each standby_name should have the form of a valid SQL identifier, unless it is *. You can use double-quoting if necessary. But note that standby_names are compared to standby application names case-insensitively, whether double-quoted or not.

If no synchronous standby names are specified here, then synchronous replication is not enabled and transaction commits will not wait for replication. This is the default configuration. Even when synchronous replication is enabled, individual transactions can be configured not to wait for replication by setting the synchronous_commit parameter to local or off.

This parameter can only be set in the postgresql.conf file or on the server command line.

`vacuum_defer_cleanup_age(integer)`

Specifies the number of transactions by which VACUUM and HOT updates will defer cleanup of dead row versions. The default is zero transactions, meaning that dead row versions can be removed as soon as possible, that is, as soon as they are no longer visible to any open transaction. You may wish to set this to a non-zero value on a primary server that is supporting hot standby servers, as described in Section 26.5. This allows more time for queries on the standby to complete without incurring conflicts due to early cleanup of rows. However, since the value is measured in terms of number of write transactions occurring on the primary server, it is difficult to predict just how much additional grace time will be made available to standby queries. This parameter can only be set in the postgresql.conf file or on the server command line.

You should also consider setting hot_standby_feedback on standby server(s) as an alternative to using this parameter.

This does not prevent cleanup of dead rows which have reached the age specified by old_s-napshot_threshold.
19.6.3. Standby Servers

These settings control the behavior of a standby server that is to receive replication data. Their values on the master server are irrelevant.

`hot_standby (boolean)`

Specifies whether or not you can connect and run queries during recovery, as described in Section 26.5. The default value is on. This parameter can only be set at server start. It only has effect during archive recovery or in standby mode.

`max_standby_archive_delay (integer)`

When Hot Standby is active, this parameter determines how long the standby server should wait before canceling standby queries that conflict with about-to-be-applied WAL entries, as described in Section 26.5.2. `max_standby_archive_delay` applies when WAL data is being read from WAL archive (and is therefore not current). The default is 30 seconds. Units are milliseconds if not specified. A value of -1 allows the standby to wait forever for conflicting queries to complete. This parameter can only be set in the `postgresql.conf` file or on the server command line.

Note that `max_standby_archive_delay` is not the same as the maximum length of time a query can run before cancellation; rather it is the maximum total time allowed to apply any one WAL segment’s data. Thus, if one query has resulted in significant delay earlier in the WAL segment, subsequent conflicting queries will have much less grace time.

`max_standby_streaming_delay (integer)`

When Hot Standby is active, this parameter determines how long the standby server should wait before canceling standby queries that conflict with about-to-be-applied WAL entries, as described in Section 26.5.2. `max_standby_streaming_delay` applies when WAL data is being received via streaming replication. The default is 30 seconds. Units are milliseconds if not specified. A value of -1 allows the standby to wait forever for conflicting queries to complete. This parameter can only be set in the `postgresql.conf` file or on the server command line.

Note that `max_standby_streaming_delay` is not the same as the maximum length of time a query can run before cancellation; rather it is the maximum total time allowed to apply WAL data once it has been received from the primary server. Thus, if one query has resulted in significant delay, subsequent conflicting queries will have much less grace time until the standby server has caught up again.

`wal_receiver_status_interval (integer)`

Specifies the minimum frequency for the WAL receiver process on the standby to send information about replication progress to the primary or upstream standby, where it can be seen using the `pg_stat_replication` view. The standby will report the last write-ahead log location it has written, the last position it has flushed to disk, and the last position it has applied. This parameter’s value is the maximum interval, in seconds, between reports. Updates are sent each time the write or flush positions change, or at least as often as specified by this parameter. Thus, the apply position may lag slightly behind the true position. Setting this parameter to zero disables status updates completely. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default value is 10 seconds.

`hot_standby_feedback (boolean)`

Specifies whether or not a hot standby will send feedback to the primary or upstream standby about queries currently executing on the standby. This parameter can be used to eliminate query cancels caused by cleanup records, but can cause database bloat on the primary for some workloads. Feedback messages will not be sent more frequently than once per `wal_receiver_s-
Server Configuration

\texttt{status\_interval}. The default value is \texttt{off}. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

If cascaded replication is in use the feedback is passed upstream until it eventually reaches the primary. Standbys make no other use of feedback they receive other than to pass upstream.

This setting does not override the behavior of \texttt{old\_snapshot\_threshold} on the primary; a snapshot on the standby which exceeds the primary's age threshold can become invalid, resulting in cancellation of transactions on the standby. This is because \texttt{old\_snapshot\_threshold} is intended to provide an absolute limit on the time which dead rows can contribute to bloat, which would otherwise be violated because of the configuration of a standby.

\texttt{wal\_receiver\_timeout (integer)}

Terminate replication connections that are inactive longer than the specified number of milliseconds. This is useful for the receiving standby server to detect a primary node crash or network outage. A value of zero disables the timeout mechanism. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line. The default value is 60 seconds.

\texttt{wal\_retrieve\_retry\_interval (integer)}

Specify how long the standby server should wait when WAL data is not available from any sources (streaming replication, local \texttt{pg\_wal} or WAL archive) before retrying to retrieve WAL data. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line. The default value is 5 seconds. Units are milliseconds if not specified.

This parameter is useful in configurations where a node in recovery needs to control the amount of time to wait for new WAL data to be available. For example, in archive recovery, it is possible to make the recovery more responsive in the detection of a new WAL log file by reducing the value of this parameter. On a system with low WAL activity, increasing it reduces the amount of requests necessary to access WAL archives, something useful for example in cloud environments where the amount of times an infrastructure is accessed is taken into account.

19.6.4. Subscribers

These settings control the behavior of a logical replication subscriber. Their values on the publisher are irrelevant.

Note that \texttt{wal\_receiver\_timeout}, \texttt{wal\_receiver\_status\_interval} and \texttt{wal\_retrieve\_retry\_interval} configuration parameters affect the logical replication workers as well.

\texttt{max\_logical\_replication\_workers (int)}

Specifies maximum number of logical replication workers. This includes both apply workers and table synchronization workers.

Logical replication workers are taken from the pool defined by \texttt{max\_worker\_processes}.

The default value is 4.

\texttt{max\_sync\_workers\_per\_subscription (integer)}

Maximum number of synchronization workers per subscription. This parameter controls the amount of parallelism of the initial data copy during the subscription initialization or when new tables are added.

Currently, there can be only one synchronization worker per table.

The synchronization workers are taken from the pool defined by \texttt{max\_logical\_replication\_workers}.
19.7. Query Planning

19.7.1. Planner Method Configuration

These configuration parameters provide a crude method of influencing the query plans chosen by the query optimizer. If the default plan chosen by the optimizer for a particular query is not optimal, a temporary solution is to use one of these configuration parameters to force the optimizer to choose a different plan. Better ways to improve the quality of the plans chosen by the optimizer include adjusting the planner cost constants (see Section 19.7.2), running ANALYZE manually, increasing the value of the default_statistics_target configuration parameter, and increasing the amount of statistics collected for specific columns using ALTER TABLE SET STATISTICS.

- **enable_bitmapscan(boolean)**
  
  Enables or disables the query planner's use of bitmap-scan plan types. The default is on.

- **enable_gathermerge(boolean)**
  
  Enables or disables the query planner's use of gather merge plan types. The default is on.

- **enable_hashagg(boolean)**
  
  Enables or disables the query planner's use of hashed aggregation plan types. The default is on.

- **enable_hashjoin(boolean)**
  
  Enables or disables the query planner's use of hash-join plan types. The default is on.

- **enable_indexscan(boolean)**
  
  Enables or disables the query planner's use of index-scan plan types. The default is on.

- **enable_indexonlyscan(boolean)**
  
  Enables or disables the query planner's use of index-only-scan plan types (see Section 11.11). The default is on.

- **enable_material(boolean)**
  
  Enables or disables the query planner's use of materialization. It is impossible to suppress materialization entirely, but turning this variable off prevents the planner from inserting materialize nodes except in cases where it is required for correctness. The default is on.

- **enable_mergejoin(boolean)**
  
  Enables or disables the query planner's use of merge-join plan types. The default is on.

- **enable_nestloop(boolean)**
  
  Enables or disables the query planner's use of nested-loop join plans. It is impossible to suppress nested-loop joins entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

- **enable_parallel_append(boolean)**
  
  Enables or disables the query planner's use of parallel-aware append plan types. The default is on.
enable_parallel_hash(boolean)

Enables or disables the query planner's use of hash-join plan types with parallel hash. Has no effect if hash-join plans are not also enabled. The default is on.

enable_partition_pruning(boolean)

Enables or disables the query planner's ability to eliminate a partitioned table's partitions from query plans. This also controls the planner's ability to generate query plans which allow the query executor to remove (ignore) partitions during query execution. The default is on. See Section 5.10.4 for details.

enable_partitionwise_join(boolean)

Enables or disables the query planner's use of partitionwise join, which allows a join between partitioned tables to be performed by joining the matching partitions. Partitionwise join currently applies only when the join conditions include all the partition keys, which must be of the same data type and have exactly matching sets of child partitions. Because partitionwise join planning can use significantly more CPU time and memory during planning, the default is off.

enable_partitionwise_aggregate(boolean)

Enables or disables the query planner's use of partitionwise grouping or aggregation, which allows grouping or aggregation on a partitioned tables performed separately for each partition. If the GROUP BY clause does not include the partition keys, only partial aggregation can be performed on a per-partition basis, and finalization must be performed later. Because partitionwise grouping or aggregation can use significantly more CPU time and memory during planning, the default is off.

enable_seqscan(boolean)

Enables or disables the query planner's use of sequential scan plan types. It is impossible to suppress sequential scans entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

enable_sort(boolean)

Enables or disables the query planner's use of explicit sort steps. It is impossible to suppress explicit sorts entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

enable_tidscan(boolean)

Enables or disables the query planner's use of TID scan plan types. The default is on.

19.7.2. Planner Cost Constants

The cost variables described in this section are measured on an arbitrary scale. Only their relative values matter, hence scaling them all up or down by the same factor will result in no change in the planner's choices. By default, these cost variables are based on the cost of sequential page fetches; that is, seq_page_cost is conventionally set to 1.0 and the other cost variables are set with reference to that. But you can use a different scale if you prefer, such as actual execution times in milliseconds on a particular machine.

Note

Unfortunately, there is no well-defined method for determining ideal values for the cost variables. They are best treated as averages over the entire mix of queries that a
particular installation will receive. This means that changing them on the basis of just a few experiments is very risky.

seq_page_cost (floating point)

Sets the planner’s estimate of the cost of a disk page fetch that is part of a series of sequential fetches. The default is 1.0. This value can be overridden for tables and indexes in a particular tablespace by setting the tablespace parameter of the same name (see ALTER TABLESPACE).

random_page_cost (floating point)

Sets the planner’s estimate of the cost of a non-sequentially-fetched disk page. The default is 4.0. This value can be overridden for tables and indexes in a particular tablespace by setting the tablespace parameter of the same name (see ALTER TABLESPACE).

Reducing this value relative to seq_page_cost will cause the system to prefer index scans; raising it will make index scans look relatively more expensive. You can raise or lower both values together to change the importance of disk I/O costs relative to CPU costs, which are described by the following parameters.

Random access to mechanical disk storage is normally much more expensive than four times sequential access. However, a lower default is used (4.0) because the majority of random accesses to disk, such as indexed reads, are assumed to be in cache. The default value can be thought of as modeling random access as 40 times slower than sequential, while expecting 90% of random reads to be cached.

If you believe a 90% cache rate is an incorrect assumption for your workload, you can increase random_page_cost to better reflect the true cost of random storage reads. Correspondingly, if your data is likely to be completely in cache, such as when the database is smaller than the total server memory, decreasing random_page_cost can be appropriate. Storage that has a low random read cost relative to sequential, e.g. solid-state drives, might also be better modeled with a lower value for random_page_cost.

Tip
Although the system will let you set random_page_cost to less than seq_page_cost, it is not physically sensible to do so. However, setting them equal makes sense if the database is entirely cached in RAM, since in that case there is no penalty for touching pages out of sequence. Also, in a heavily-cached database you should lower both values relative to the CPU parameters, since the cost of fetching a page already in RAM is much smaller than it would normally be.

cpu_tuple_cost (floating point)

Sets the planner’s estimate of the cost of processing each row during a query. The default is 0.01.

cpu_index_tuple_cost (floating point)

Sets the planner’s estimate of the cost of processing each index entry during an index scan. The default is 0.005.

cpu_operator_cost (floating point)

Sets the planner’s estimate of the cost of processing each operator or function executed during a query. The default is 0.0025.
Server Configuration

parallel_setup_cost (floating point)

Sets the planner's estimate of the cost of launching parallel worker processes. The default is 1000.

parallel_tuple_cost (floating point)

Sets the planner's estimate of the cost of transferring one tuple from a parallel worker process to another process. The default is 0.1.

min_parallel_table_scan_size (integer)

Sets the minimum amount of table data that must be scanned in order for a parallel scan to be considered. For a parallel sequential scan, the amount of table data scanned is always equal to the size of the table, but when indexes are used the amount of table data scanned will normally be less. The default is 8 megabytes (8MB).

min_parallel_index_scan_size (integer)

Sets the minimum amount of index data that must be scanned in order for a parallel scan to be considered. Note that a parallel index scan typically won't touch the entire index; it is the number of pages which the planner believes will actually be touched by the scan which is relevant. The default is 512 kilobytes (512kB).

effective_cache_size (integer)

Sets the planner's assumption about the effective size of the disk cache that is available to a single query. This is factored into estimates of the cost of using an index; a higher value makes it more likely index scans will be used, a lower value makes it more likely sequential scans will be used. When setting this parameter you should consider both PostgreSQL's shared buffers and the portion of the kernel's disk cache that will be used for PostgreSQL data files, though some data might exist in both places. Also, take into account the expected number of concurrent queries on different tables, since they will have to share the available space. This parameter has no effect on the size of shared memory allocated by PostgreSQL, nor does it reserve kernel disk cache; it is used only for estimation purposes. The system also does not assume data remains in the disk cache between queries. The default is 4 gigabytes (4GB).

jit_above_cost (floating point)

Sets the query cost above which JIT compilation is activated, if enabled (see Chapter 32). Performing JIT costs planning time but can accelerate query execution. Setting this to -1 disables JIT compilation. The default is 100000.

jit_inline_above_cost (floating point)

Sets the query cost above which JIT compilation attempts to inline functions and operators. Inlining adds planning time, but can improve execution speed. It is not meaningful to set this to less than jit_above_cost. Setting this to -1 disables inlining. The default is 500000.

jit_optimize_above_cost (floating point)

Sets the query cost above which JIT compilation applies expensive optimizations. Such optimization adds planning time, but can improve execution speed. It is not meaningful to set this to less than jit_above_cost, and it is unlikely to be beneficial to set it to more than jit_inline_above_cost. Setting this to -1 disables expensive optimizations. The default is 500000.

19.7.3. Genetic Query Optimizer

The genetic query optimizer (GEQO) is an algorithm that does query planning using heuristic searching. This reduces planning time for complex queries (those joining many relations), at the cost of pro-
ducing plans that are sometimes inferior to those found by the normal exhaustive-search algorithm. For more information see Chapter 60.

\textbf{geqo (boolean)}

Enables or disables genetic query optimization. This is on by default. It is usually best not to turn it off in production; the \texttt{geqo\_threshold} variable provides more granular control of GEQO.

\textbf{geqo\_threshold (integer)}

Use genetic query optimization to plan queries with at least this many \texttt{FROM} items involved. (Note that a \texttt{FULL OUTER JOIN} construct counts as only one \texttt{FROM} item.) The default is 12. For simpler queries it is usually best to use the regular, exhaustive-search planner, but for queries with many tables the exhaustive search takes too long, often longer than the penalty of executing a suboptimal plan. Thus, a threshold on the size of the query is a convenient way to manage use of GEQO.

\textbf{geqo\_effort (integer)}

Controls the trade-off between planning time and query plan quality in GEQO. This variable must be an integer in the range from 1 to 10. The default value is five. Larger values increase the time spent doing query planning, but also increase the likelihood that an efficient query plan will be chosen.

\texttt{geqo\_effort} doesn't actually do anything directly; it is only used to compute the default values for the other variables that influence GEQO behavior (described below). If you prefer, you can set the other parameters by hand instead.

\textbf{geqo\_pool\_size (integer)}

Controls the pool size used by GEQO, that is the number of individuals in the genetic population. It must be at least two, and useful values are typically 100 to 1000. If it is set to zero (the default setting) then a suitable value is chosen based on \texttt{geqo\_effort} and the number of tables in the query.

\textbf{geqo\_generations (integer)}

Controls the number of generations used by GEQO, that is the number of iterations of the algorithm. It must be at least one, and useful values are in the same range as the pool size. If it is set to zero (the default setting) then a suitable value is chosen based on \texttt{geqo\_pool\_size}.

\textbf{geqo\_selection\_bias (floating point)}

Controls the selection bias used by GEQO. The selection bias is the selective pressure within the population. Values can be from 1.50 to 2.00; the latter is the default.

\textbf{geqo\_seed (floating point)}

Controls the initial value of the random number generator used by GEQO to select random paths through the join order search space. The value can range from zero (the default) to one. Varying the value changes the set of join paths explored, and may result in a better or worse best path being found.

\section*{19.7.4. Other Planner Options}

\textbf{default\_statistics\_target (integer)}

Sets the default statistics target for table columns without a column-specific target set via \texttt{ALTER TABLE SET STATISTICS}. Larger values increase the time needed to do \texttt{ANALYZE}, but might
improve the quality of the planner's estimates. The default is 100. For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.

**constraint_exclusion (enum)**

Controls the query planner's use of table constraints to optimize queries. The allowed values of constraint_exclusion are on (examine constraints for all tables), off (never examine constraints), and partition (examine constraints only for inheritance child tables and UNION ALL subqueries). partition is the default setting. It is often used with inheritance tables to improve performance.

When this parameter allows it for a particular table, the planner compares query conditions with the table's CHECK constraints, and omits scanning tables for which the conditions contradict the constraints. For example:

```sql
CREATE TABLE parent(key integer, ...);
CREATE TABLE child1000(check (key between 1000 and 1999)) INHERITS(parent);
CREATE TABLE child2000(check (key between 2000 and 2999)) INHERITS(parent);
...
SELECT * FROM parent WHERE key = 2400;
```

With constraint exclusion enabled, this `SELECT` will not scan `child1000` at all, improving performance.

Currently, constraint exclusion is enabled by default only for cases that are often used to implement table partitioning via inheritance tables. Turning it on for all tables imposes extra planning overhead that is quite noticeable on simple queries, and most often will yield no benefit for simple queries. If you have no inheritance partitioned tables you might prefer to turn it off entirely.

Refer to Section 5.10.5 for more information on using constraint exclusion and partitioning.

**cursor_tuple_fraction (floating point)**

Sets the planner's estimate of the fraction of a cursor's rows that will be retrieved. The default is 0.1. Smaller values of this setting bias the planner towards using “fast start” plans for cursors, which will retrieve the first few rows quickly while perhaps taking a long time to fetch all rows. Larger values put more emphasis on the total estimated time. At the maximum setting of 1.0, cursors are planned exactly like regular queries, considering only the total estimated time and not how soon the first rows might be delivered.

**fromCollapseLimit (integer)**

The planner will merge sub-queries into upper queries if the resulting FROM list would have no more than this many items. Smaller values reduce planning time but might yield inferior query plans. The default is eight. For more information see Section 14.3.

Setting this value to geqo_threshold or more may trigger use of the GEQO planner, resulting in non-optimal plans. See Section 19.7.3.

**jit (boolean)**

Determines whether JIT compilation may be used by PostgreSQL, if available (see Chapter 32). The default is off.

**joinCollapseLimit (integer)**

The planner will rewrite explicit JOIN constructs (except FULL JOINs) into lists of FROM items whenever a list of no more than this many items would result. Smaller values reduce planning time but might yield inferior query plans.
By default, this variable is set the same as `fromCollapseLimit`, which is appropriate for most uses. Setting it to 1 prevents any reordering of explicit `JOIN`s. Thus, the explicit join order specified in the query will be the actual order in which the relations are joined. Because the query planner does not always choose the optimal join order, advanced users can elect to temporarily set this variable to 1, and then specify the join order they desire explicitly. For more information see Section 14.3.

Setting this value to `geqoThreshold` or more may trigger use of the GEQO planner, resulting in non-optimal plans. See Section 19.7.3.

**parallel_leader_participation (boolean)**

Allows the leader process to execute the query plan under `Gather` and `Gather Merge` nodes instead of waiting for worker processes. The default is `on`. Setting this value to `off` reduces the likelihood that workers will become blocked because the leader is not reading tuples fast enough, but requires the leader process to wait for worker processes to start up before the first tuples can be produced. The degree to which the leader can help or hinder performance depends on the plan type, number of workers and query duration.

**force_parallel_mode (enum)**

Allows the use of parallel queries for testing purposes even in cases where no performance benefit is expected. The allowed values of `force_parallel_mode` are `off` (use parallel mode only when it is expected to improve performance), `on` (force parallel query for all queries for which it is thought to be safe), and `regress` (like `on`, but with additional behavior changes as explained below).

More specifically, setting this value to `on` will add a `Gather` node to the top of any query plan for which this appears to be safe, so that the query runs inside of a parallel worker. Even when a parallel worker is not available or cannot be used, operations such as starting a subtransaction that would be prohibited in a parallel query context will be prohibited unless the planner believes that this will cause the query to fail. If failures or unexpected results occur when this option is set, some functions used by the query may need to be marked `PARALLEL UNSAFE` (or, possibly, `PARALLEL RESTRICTED`).

Setting this value to `regress` has all of the same effects as setting it to `on` plus some additional effects that are intended to facilitate automated regression testing. Normally, messages from a parallel worker include a context line indicating that, but a setting of `regress` suppresses this line so that the output is the same as in non-parallel execution. Also, the `Gather` nodes added to plans by this setting are hidden in `EXPLAIN` output so that the output matches what would be obtained if this setting were turned `off`.

### 19.8. Error Reporting and Logging

#### 19.8.1. Where To Log

**log_destination (string)**

PostgreSQL supports several methods for logging server messages, including `stderr`, `csvlog` and `syslog`. On Windows, `eventlog` is also supported. Set this parameter to a list of desired log destinations separated by commas. The default is to log to `stderr` only. This parameter can only be set in the `postgresql.conf` file or on the server command line.

If `csvlog` is included in `log_destination`, log entries are output in “comma separated value” (CSV) format, which is convenient for loading logs into programs. See Section 19.8.4 for details. `logging_collector` must be enabled to generate CSV-format log output.

When either `stderr` or `csvlog` are included, the file `current_logfiles` is created to record the location of the log file(s) currently in use by the logging collector and the associated logging
destination. This provides a convenient way to find the logs currently in use by the instance. Here is an example of this file's content:

stderr log/postgresql.log
csvlog log/postgresql.csv

current_logfiles is recreated when a new log file is created as an effect of rotation, and when log_destination is reloaded. It is removed when neither stderr nor csvlog are included in log_destination, and when the logging collector is disabled.

Note

On most Unix systems, you will need to alter the configuration of your system's syslog daemon in order to make use of the syslog option for log_destination. PostgreSQL can log to syslog facilities LOCAL0 through LOCAL7 (see syslog_facility), but the default syslog configuration on most platforms will discard all such messages. You will need to add something like:

local0.* /var/log/postgresql

to the syslog daemon's configuration file to make it work.

On Windows, when you use the eventlog option for log_destination, you should register an event source and its library with the operating system so that the Windows Event Viewer can display event log messages cleanly. See Section 18.11 for details.

logging_collector (boolean)

This parameter enables the logging collector, which is a background process that captures log messages sent to stderr and redirects them into log files. This approach is often more useful than logging to syslog, since some types of messages might not appear in syslog output. (One common example is dynamic-linker failure messages; another is error messages produced by scripts such as archive_command.) This parameter can only be set at server start.

Note

It is possible to log to stderr without using the logging collector; the log messages will just go to wherever the server's stderr is directed. However, that method is only suitable for low log volumes, since it provides no convenient way to rotate log files. Also, on some platforms not using the logging collector can result in lost or garbled log output, because multiple processes writing concurrently to the same log file can overwrite each other's output.

Note

The logging collector is designed to never lose messages. This means that in case of extremely high load, server processes could be blocked while trying to send additional log messages when the collector has fallen behind. In contrast, syslog prefers to drop messages if it cannot write them, which means it may fail to log some messages in such cases but it will not block the rest of the system.
log_directory (string)

When logging_collector is enabled, this parameter determines the directory in which log files will be created. It can be specified as an absolute path, or relative to the cluster data directory. This parameter can only be set in the postgresql.conf file or on the server command line. The default is log.

log_filename (string)

When logging_collector is enabled, this parameter sets the file names of the created log files. The value is treated as a strftime pattern, so % escapes can be used to specify time-varying file names. (Note that if there are any time-zone-dependent %-escapes, the computation is done in the zone specified by log_timezone.) The supported %-escapes are similar to those listed in the Open Group’s strftime specification. Note that the system’s strftime is not used directly, so platform-specific (nonstandard) extensions do not work. The default is postgresql-%Y-%m-%d_%H%M%S.log.

If you specify a file name without escapes, you should plan to use a log rotation utility to avoid eventually filling the entire disk. In releases prior to 8.4, if no % escapes were present, PostgreSQL would append the epoch of the new log file’s creation time, but this is no longer the case.

If CSV-format output is enabled in log_destination, .csv will be appended to the time-stamped log file name to create the file name for CSV-format output. (If log_filename ends in .log, the suffix is replaced instead.)

This parameter can only be set in the postgresql.conf file or on the server command line.

log_file_mode (integer)

On Unix systems this parameter sets the permissions for log files when logging_collector is enabled. (On Microsoft Windows this parameter is ignored.) The parameter value is expected to be a numeric mode specified in the format accepted by the chmod and umask system calls. (To use the customary octal format the number must start with a 0 (zero).)

The default permissions are 0600, meaning only the server owner can read or write the log files. The other commonly useful setting is 0640, allowing members of the owner's group to read the files. Note however that to make use of such a setting, you'll need to alter log_directory to store the files somewhere outside the cluster data directory. In any case, it’s unwise to make the log files world-readable, since they might contain sensitive data.

This parameter can only be set in the postgresql.conf file or on the server command line.

log_rotation_age (integer)

When logging_collector is enabled, this parameter determines the maximum lifetime of an individual log file. After this many minutes have elapsed, a new log file will be created. Set to zero to disable time-based creation of new log files. This parameter can only be set in the postgresql.conf file or on the server command line.

log_rotation_size (integer)

When logging_collector is enabled, this parameter determines the maximum size of an individual log file. After this many kilobytes have been emitted into a log file, a new log file will be created. Set to zero to disable size-based creation of new log files. This parameter can only be set in the postgresql.conf file or on the server command line.

log_truncate_on_rotation (boolean)

When logging_collector is enabled, this parameter will cause PostgreSQL to truncate (overwrite), rather than append to, any existing log file of the same name. However, truncation
will occur only when a new file is being opened due to time-based rotation, not during server startup or size-based rotation. When off, pre-existing files will be appended to in all cases. For example, using this setting in combination with a log_filename like postgresql-%H.log would result in generating twenty-four hourly log files and then cyclically overwriting them. This parameter can only be set in the postgresql.conf file or on the server command line.

Example: To keep 7 days of logs, one log file per day named server_log.Mon, server_log.Tue, etc, and automatically overwrite last week's log with this week's log, set log_filename to server_log.%a, log_truncate_on_rotation to on, and log_rotation_age to 1440.

Example: To keep 24 hours of logs, one log file per hour, but also rotate sooner if the log file size exceeds 1GB, set log_filename to server_log.%H%M, log_truncate_on_rotation to on, log_rotation_age to 60, and log_rotation_size to 1000000. Including %M in log_filename allows any size-driven rotations that might occur to select a file name different from the hour's initial file name.

syslog_facility(enum)

When logging to syslog is enabled, this parameter determines the syslog “facility” to be used. You can choose from LOCAL0, LOCAL1, LOCAL2, LOCAL3, LOCAL4, LOCAL5, LOCAL6, LOCAL7; the default is LOCAL0. See also the documentation of your system's syslog daemon.

This parameter can only be set in the postgresql.conf file or on the server command line.

syslog_ident(string)

When logging to syslog is enabled, this parameter determines the program name used to identify PostgreSQL messages in syslog logs. The default is postgres. This parameter can only be set in the postgresql.conf file or on the server command line.

syslog_sequence_numbers(boolean)

When logging to syslog and this is on (the default), then each message will be prefixed by an increasing sequence number (such as [2]). This circumvents the “--- last message repeated N times ---” suppression that many syslog implementations perform by default. In more modern syslog implementations, repeated message suppression can be configured (for example, $RepeatedMsgReduction in rsyslog), so this might not be necessary. Also, you could turn this off if you actually want to suppress repeated messages.

This parameter can only be set in the postgresql.conf file or on the server command line.

syslog_split_messages(boolean)

When logging to syslog is enabled, this parameter determines how messages are delivered to syslog. When on (the default), messages are split by lines, and long lines are split so that they will fit into 1024 bytes, which is a typical size limit for traditional syslog implementations. When off, PostgreSQL server log messages are delivered to the syslog service as is, and it is up to the syslog service to cope with the potentially bulky messages.

If syslog is ultimately logging to a text file, then the effect will be the same either way, and it is best to leave the setting on, since most syslog implementations either cannot handle large messages or would need to be specially configured to handle them. But if syslog is ultimately writing into some other medium, it might be necessary or more useful to keep messages logically together.

This parameter can only be set in the postgresql.conf file or on the server command line.

event_source(string)

When logging to event log is enabled, this parameter determines the program name used to identify PostgreSQL messages in the log. The default is PostgreSQL. This parameter can only be set in the postgresql.conf file or on the server command line.
19.8.2. When To Log

client_min_messages (enum)

Controls which message levels are sent to the client. Valid values are DEBUG5, DEBUG4, DEBUG3, DEBUG2, DEBUG1, LOG, NOTICE, WARNING, ERROR, FATAL, and PANIC. Each level includes all the levels that follow it. The later the level, the fewer messages are sent. The default is NOTICE. Note that LOG has a different rank here than in log_min_messages.

log_min_messages (enum)

Controls which message levels are written to the server log. Valid values are DEBUG5, DEBUG4, DEBUG3, DEBUG2, DEBUG1, INFO, NOTICE, WARNING, ERROR, LOG, FATAL, and PANIC. Each level includes all the levels that follow it. The later the level, the fewer messages are sent to the log. The default is WARNING. Note that LOG has a different rank here than in client_min_messages. Only superusers can change this setting.

log_min_error_statement (enum)

Controls which SQL statements that cause an error condition are recorded in the server log. The current SQL statement is included in the log entry for any message of the specified severity or higher. Valid values are DEBUG5, DEBUG4, DEBUG3, DEBUG2, DEBUG1, INFO, NOTICE, WARNING, ERROR, LOG, FATAL, and PANIC. The default is ERROR, which means statements causing errors, log messages, fatal errors, or panics will be logged. To effectively turn off logging of failing statements, set this parameter to PANIC. Only superusers can change this setting.

log_min_duration_statement (integer)

Causes the duration of each completed statement to be logged if the statement ran for at least the specified number of milliseconds. Setting this to zero prints all statement durations. Minus-one (the default) disables logging statement durations. For example, if you set it to 250ms then all SQL statements that run 250ms or longer will be logged. Enabling this parameter can be helpful in tracking down unoptimized queries in your applications. Only superusers can change this setting.

For clients using extended query protocol, durations of the Parse, Bind, and Execute steps are logged independently.

**Note**

When using this option together with log_statement, the text of statements that are logged because of log_statement will not be repeated in the duration log message. If you are not using syslog, it is recommended that you log the PID or session ID using log_line_prefix so that you can link the statement message to the later duration message using the process ID or session ID.

Table 19.1 explains the message severity levels used by PostgreSQL. If logging output is sent to syslog or Windows' eventlog, the severity levels are translated as shown in the table.

**Table 19.1. Message Severity Levels**

<table>
<thead>
<tr>
<th>Severity</th>
<th>Usage</th>
<th>syslog</th>
<th>eventlog</th>
</tr>
</thead>
<tbody>
<tr>
<td>DEBUG1..DEBUG5</td>
<td>Provides successively-more-detailed information for use by developers.</td>
<td>DEBUG</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>Severity</td>
<td>Usage</td>
<td>syslog</td>
<td>eventlog</td>
</tr>
<tr>
<td>----------</td>
<td>-------</td>
<td>--------</td>
<td>----------</td>
</tr>
<tr>
<td>INFO</td>
<td>Provides information implicitly requested by the user, e.g., output from VACUUM VERBOSE.</td>
<td>INFO</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>NOTICE</td>
<td>Provides information that might be helpful to users, e.g., notice of truncation of long identifiers.</td>
<td>NOTICE</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>WARNING</td>
<td>Provides warnings of likely problems, e.g., COMMIT outside a transaction block.</td>
<td>NOTICE</td>
<td>WARNING</td>
</tr>
<tr>
<td>ERROR</td>
<td>Reports an error that caused the current command to abort.</td>
<td>WARNING</td>
<td>ERROR</td>
</tr>
<tr>
<td>LOG</td>
<td>Reports information of interest to administrators, e.g., checkpoint activity.</td>
<td>INFO</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>FATAL</td>
<td>Reports an error that caused the current session to abort.</td>
<td>ERR</td>
<td>ERROR</td>
</tr>
<tr>
<td>PANIC</td>
<td>Reports an error that caused all database sessions to abort.</td>
<td>CRIT</td>
<td>ERROR</td>
</tr>
</tbody>
</table>

### 19.8.3. What To Log

**application_name (string)**

The application_name can be any string of less than NAMEDATALEN characters (64 characters in a standard build). It is typically set by an application upon connection to the server. The name will be displayed in the pg_stat_activity view and included in CSV log entries. It can also be included in regular log entries via the log_line_prefix parameter. Only printable ASCII characters may be used in the application_name value. Other characters will be replaced with question marks (?).

**debug_print_parse (boolean)**
**debug_print_rewritten (boolean)**
**debug_print_plan (boolean)**

These parameters enable various debugging output to be emitted. When set, they print the resulting parse tree, the query rewriter output, or the execution plan for each executed query. These messages are emitted at LOG message level, so by default they will appear in the server log but will not be sent to the client. You can change that by adjusting client_min_messages and/or log_min_messages. These parameters are off by default.

**debug_pretty_print (boolean)**

When set, debug_pretty_print indents the messages produced by debug_print_parse, debug_print_rewritten, or debug_print_plan. This results in more readable but much longer output than the "compact" format used when it is off. It is on by default.
log_checkpoints (boolean)

Causes checkpoints and restartpoints to be logged in the server log. Some statistics are included in the log messages, including the number of buffers written and the time spent writing them. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is off.

log_connections (boolean)

Causes each attempted connection to the server to be logged, as well as successful completion of client authentication. Only superusers can change this parameter at session start, and it cannot be changed at all within a session. The default is off.

**Note**

Some client programs, like `psql`, attempt to connect twice while determining if a password is required, so duplicate “connection received” messages do not necessarily indicate a problem.

log_disconnections (boolean)

Causes session terminations to be logged. The log output provides information similar to `log_connections`, plus the duration of the session. Only superusers can change this parameter at session start, and it cannot be changed at all within a session. The default is off.

log_duration (boolean)

Causes the duration of every completed statement to be logged. The default is off. Only superusers can change this setting.

For clients using extended query protocol, durations of the Parse, Bind, and Execute steps are logged independently.

**Note**

The difference between setting this option and setting `log_min_duration_statement` to zero is that exceeding `log_min_duration_statement` forces the text of the query to be logged, but this option doesn't. Thus, if `log_duration` is on and `log_min_duration_statement` has a positive value, all durations are logged but the query text is included only for statements exceeding the threshold. This behavior can be useful for gathering statistics in high-load installations.

log_error_verbosity (enum)

Controls the amount of detail written in the server log for each message that is logged. Valid values are TERSE, DEFAULT, and VERBOSE, each adding more fields to displayed messages. TERSE excludes the logging of DETAIL, HINT, QUERY, and CONTEXT error information. VERBOSE output includes the SQLSTATE error code (see also Appendix A) and the source code file name, function name, and line number that generated the error. Only superusers can change this setting.

log_hostname (boolean)

By default, connection log messages only show the IP address of the connecting host. Turning this parameter on causes logging of the host name as well. Note that depending on your host name
resolution setup this might impose a non-negligible performance penalty. This parameter can only be set in the `postgresql.conf` file or on the server command line.

**log_line_prefix**(string)

This is a printf-style string that is output at the beginning of each log line. % characters begin "escape sequences" that are replaced with status information as outlined below. Unrecognized escapes are ignored. Other characters are copied straight to the log line. Some escapes are only recognized by session processes, and will be treated as empty by background processes such as the main server process. Status information may be aligned either left or right by specifying a numeric literal after the % and before the option. A negative value will cause the status information to be padded on the right with spaces to give it a minimum width, whereas a positive value will pad on the left. Padding can be useful to aid human readability in log files. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is '%m [%p]' which logs a time stamp and the process ID.

<table>
<thead>
<tr>
<th>Escape</th>
<th>Effect</th>
<th>Session only</th>
</tr>
</thead>
<tbody>
<tr>
<td>%a</td>
<td>Application name</td>
<td>yes</td>
</tr>
<tr>
<td>%u</td>
<td>User name</td>
<td>yes</td>
</tr>
<tr>
<td>%d</td>
<td>Database name</td>
<td>yes</td>
</tr>
<tr>
<td>%r</td>
<td>Remote host name or IP address, and remote port</td>
<td>yes</td>
</tr>
<tr>
<td>%h</td>
<td>Remote host name or IP address</td>
<td>yes</td>
</tr>
<tr>
<td>%p</td>
<td>Process ID</td>
<td>no</td>
</tr>
<tr>
<td>%t</td>
<td>Time stamp without milliseconds</td>
<td>no</td>
</tr>
<tr>
<td>%m</td>
<td>Time stamp with milliseconds</td>
<td>no</td>
</tr>
<tr>
<td>%n</td>
<td>Time stamp with milliseconds (as a Unix epoch)</td>
<td>no</td>
</tr>
<tr>
<td>%i</td>
<td>Command tag: type of session's current command</td>
<td>yes</td>
</tr>
<tr>
<td>%e</td>
<td>SQLSTATE error code</td>
<td>no</td>
</tr>
<tr>
<td>%c</td>
<td>Session ID: see below</td>
<td>no</td>
</tr>
<tr>
<td>%l</td>
<td>Number of the log line for each session or process, starting at 1</td>
<td>no</td>
</tr>
<tr>
<td>%s</td>
<td>Process start time stamp</td>
<td>no</td>
</tr>
<tr>
<td>%v</td>
<td>Virtual transaction ID (backendID/localXID)</td>
<td>no</td>
</tr>
<tr>
<td>%x</td>
<td>Transaction ID (0 if none is assigned)</td>
<td>no</td>
</tr>
<tr>
<td>%q</td>
<td>Produces no output, but tells non-session processes to stop at this point in the string; ignored by session processes</td>
<td>no</td>
</tr>
<tr>
<td>%%</td>
<td>Literal %</td>
<td>no</td>
</tr>
</tbody>
</table>

The %c escape prints a quasi-unique session identifier, consisting of two 4-byte hexadecimal numbers (without leading zeros) separated by a dot. The numbers are the process start time and the process ID, so %c can also be used as a space saving way of printing those items. For example, to generate the session identifier from `pg_stat_activity`, use this query:
Server Configuration

```sql
SELECT to_hex(trunc(EXTRACT(EPOCH FROM backend_start))::integer)
   || '.' || to_hex(pid)
FROM pg_stat_activity;
```

**Tip**

If you set a nonempty value for `log_line_prefix`, you should usually make its last character be a space, to provide visual separation from the rest of the log line. A punctuation character can be used too.

**Tip**

Syslog produces its own time stamp and process ID information, so you probably do not want to include those escapes if you are logging to syslog.

**Tip**

The `%q` escape is useful when including information that is only available in session (backend) context like user or database name. For example:

```sql
log_line_prefix = '%m [%p] %q%u@%d/%a'
```

---

**log_lock_waits** (boolean)

Controls whether a log message is produced when a session waits longer than `deadlock_timeout` to acquire a lock. This is useful in determining if lock waits are causing poor performance. The default is *off*. Only superusers can change this setting.

**log_statement** (enum)

Controls which SQL statements are logged. Valid values are `none` (*off*), `ddl`, `mod`, and `all` (all statements). `ddl` logs all data definition statements, such as `CREATE`, `ALTER`, and `DROP` statements. `mod` logs all `ddl` statements, plus data-modifying statements such as `INSERT`, `UPDATE`, `DELETE`, `TRUNCATE`, and `COPY FROM PREPARE`, `EXECUTE`, and `EXPLAIN ANALYZE` statements are also logged if their contained command is of an appropriate type. For clients using extended query protocol, logging occurs when an `Execute` message is received, and values of the `Bind` parameters are included (with any embedded single-quote marks doubled).

The default is *none*. Only superusers can change this setting.

**Note**

Statements that contain simple syntax errors are not logged even by the `log_statement = all` setting, because the log message is emitted only after basic parsing has been done to determine the statement type. In the case of extended query protocol, this setting likewise does not log statements that fail before the `Execute` phase (i.e., during parse analysis or planning). Set `log_min_error_statement` to `ERROR` (or lower) to log such statements.
Server Configuration

log_replication_commands (boolean)

Causes each replication command to be logged in the server log. See Section 53.4 for more information about replication command. The default value is off. Only superusers can change this setting.

log_temp_files (integer)

Controls logging of temporary file names and sizes. Temporary files can be created for sorts, hashes, and temporary query results. A log entry is made for each temporary file when it is deleted. A value of zero logs all temporary file information, while positive values log only files whose size is greater than or equal to the specified number of kilobytes. The default setting is -1, which disables such logging. Only superusers can change this setting.

log_timezone (string)

Sets the time zone used for timestamps written in the server log. Unlike TimeZone, this value is cluster-wide, so that all sessions will report timestamps consistently. The built-in default is GMT, but that is typically overridden in postgresql.conf; initdb will install a setting there corresponding to its system environment. See Section 8.5.3 for more information. This parameter can only be set in the postgresql.conf file or on the server command line.

19.8.4. Using CSV-Format Log Output

Including csvlog in the log_destination list provides a convenient way to import log files into a database table. This option emits log lines in comma-separated-values (CSV) format, with these columns: time stamp with milliseconds, user name, database name, process ID, client host:port number, session ID, per-session line number, command tag, session start time, virtual transaction ID, regular transaction ID, error severity, SQLSTATE code, error message, error message detail, hint, internal query that led to the error (if any), character count of the error position therein, error context, user query that led to the error (if any and enabled by log_min_error_statement), character count of the error position therein, location of the error in the PostgreSQL source code (if log_error_verbosity is set to verbose), and application name. Here is a sample table definition for storing CSV-format log output:

```sql
CREATE TABLE postgres_log
(
  log_time timestamp(3) with time zone,
  user_name text,
  database_name text,
  process_id integer,
  connection_from text,
  session_id text,
  session_line_num bigint,
  command_tag text,
  session_start_time timestamp with time zone,
  virtual_transaction_id text,
  transaction_id bigint,
  error_severity text,
  sql_state_code text,
  message text,
  detail text,
  hint text,
  internal_query text,
  internal_query_pos integer,
  context text,
  query text,
)"
To import a log file into this table, use the COPY FROM command:

COPY postgres_log FROM '/full/path/to/logfile.csv' WITH csv;

There are a few things you need to do to simplify importing CSV log files:

1. Set log_filename and log_rotation_age to provide a consistent, predictable naming scheme for your log files. This lets you predict what the file name will be and know when an individual log file is complete and therefore ready to be imported.

2. Set log_rotation_size to 0 to disable size-based log rotation, as it makes the log file name difficult to predict.

3. Set log_truncate_on_rotation to on so that old log data isn't mixed with the new in the same file.

4. The table definition above includes a primary key specification. This is useful to protect against accidentally importing the same information twice. The COPY command commits all of the data it imports at one time, so any error will cause the entire import to fail. If you import a partial log file and later import the file again when it is complete, the primary key violation will cause the import to fail. Wait until the log is complete and closed before importing. This procedure will also protect against accidentally importing a partial line that hasn't been completely written, which would also cause COPY to fail.

### 19.8.5. Process Title

These settings control how process titles of server processes are modified. Process titles are typically viewed using programs like ps or, on Windows, Process Explorer. See Section 28.1 for details.

**cluster_name (string)**

Sets the cluster name that appears in the process title for all server processes in this cluster. The name can be any string of less than NAMEDATALEN characters (64 characters in a standard build). Only printable ASCII characters may be used in the cluster_name value. Other characters will be replaced with question marks (?). No name is shown if this parameter is set to the empty string ' ' (which is the default). This parameter can only be set at server start.

**update_process_title (boolean)**

Enables updating of the process title every time a new SQL command is received by the server. This setting defaults to on on most platforms, but it defaults to off on Windows due to that platform's larger overhead for updating the process title. Only superusers can change this setting.

### 19.9. Run-time Statistics

#### 19.9.1. Query and Index Statistics Collector

These parameters control server-wide statistics collection features. When statistics collection is enabled, the data that is produced can be accessed via the pg_stat and pg_statio family of system views. Refer to Chapter 28 for more information.
track_activities (boolean)

Enables the collection of information on the currently executing command of each session, along with the time when that command began execution. This parameter is on by default. Note that even when enabled, this information is not visible to all users, only to superusers and the user owning the session being reported on, so it should not represent a security risk. Only superusers can change this setting.

track_activity_query_size (integer)

Specifies the number of bytes reserved to track the currently executing command for each active session, for the pg_stat_activity.query field. The default value is 1024. This parameter can only be set at server start.

track_counts (boolean)

Enables collection of statistics on database activity. This parameter is on by default, because the autovacuum daemon needs the collected information. Only superusers can change this setting.

track_io_timing (boolean)

Enables timing of database I/O calls. This parameter is off by default, because it will repeatedly query the operating system for the current time, which may cause significant overhead on some platforms. You can use the pg_test_timing tool to measure the overhead of timing on your system. I/O timing information is displayed in pg_stat_database, in the output of EXPLAIN when the BUFFERS option is used, and by pg_stat_statements. Only superusers can change this setting.

track_functions (enum)

Enables tracking of function call counts and time used. Specify pl to track only procedural-language functions, all to also track SQL and C language functions. The default is none, which disables function statistics tracking. Only superusers can change this setting.

Note

SQL-language functions that are simple enough to be “inlined” into the calling query will not be tracked, regardless of this setting.

stats_temp_directory (string)

Sets the directory to store temporary statistics data in. This can be a path relative to the data directory or an absolute path. The default is pg_stat.tmp. Pointing this at a RAM-based file system will decrease physical I/O requirements and can lead to improved performance. This parameter can only be set in the postgresql.conf file or on the server command line.

19.9.2. Statistics Monitoring

log_statement_stats (boolean)
log_parser_stats (boolean)
log_planner_stats (boolean)
log_executor_stats (boolean)

For each query, output performance statistics of the respective module to the server log. This is a crude profiling instrument, similar to the Unix getrusage() operating system facility. log_statement_stats reports total statement statistics, while the others report per-module statistics. log_statement_stats cannot be enabled together with any of the per-module options. All of these options are disabled by default. Only superusers can change these settings.
19.10. Automatic Vacuuming

These settings control the behavior of the *autovacuum* feature. Refer to Section 24.1.6 for more information. Note that many of these settings can be overridden on a per-table basis; see Storage Parameters.

autovacuum (boolean)

Controls whether the server should run the autovacuum launcher daemon. This is on by default; however, track_counts must also be enabled for autovacuum to work. This parameter can only be set in the `postgresql.conf` file or on the server command line; however, autovacuuming can be disabled for individual tables by changing table storage parameters.

Note that even when this parameter is disabled, the system will launch autovacuum processes if necessary to prevent transaction ID wraparound. See Section 24.1.5 for more information.

log_autovacuum_min_duration (integer)

Causes each action executed by autovacuum to be logged if it ran for at least the specified number of milliseconds. Setting this to zero logs all autovacuum actions. Minus-one (the default) disables logging autovacuum actions. For example, if you set this to 250ms then all automatic vacuums and analyzes that run 250ms or longer will be logged. In addition, when this parameter is set to any value other than -1, a message will be logged if an autovacuum action is skipped due to a conflicting lock or a concurrently dropped relation. Enabling this parameter can be helpful in tracking autovacuum activity. This parameter can only be set in the `postgresql.conf` file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_max_workers (integer)

Specifies the maximum number of autovacuum processes (other than the autovacuum launcher) that may be running at any one time. The default is three. This parameter can only be set at server start.

autovacuum_naptime (integer)

Specifies the minimum delay between autovacuum runs on any given database. In each round the daemon examines the database and issues VACUUM and ANALYZE commands as needed for tables in that database. The delay is measured in seconds, and the default is one minute (1min). This parameter can only be set in the `postgresql.conf` file or on the server command line.

autovacuum_vacuum_threshold (integer)

Specifies the minimum number of updated or deleted tuples needed to trigger a VACUUM in any one table. The default is 50 tuples. This parameter can only be set in the `postgresql.conf` file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_analyze_threshold (integer)

Specifies the minimum number of inserted, updated or deleted tuples needed to trigger an ANALYZE in any one table. The default is 50 tuples. This parameter can only be set in the `postgresql.conf` file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_vacuum_scale_factor (floating point)

Specifies a fraction of the table size to add to `autovacuum_vacuum_threshold` when deciding whether to trigger a VACUUM. The default is 0.2 (20% of table size). This parameter can
Server Configuration

only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_analyze_scale_factor(floating point)

Specifies a fraction of the table size to add to autovacuum_analyze_threshold when deciding whether to trigger an ANALYZE. The default is 0.1 (10% of table size). This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_freeze_max_age(integer)

Specifies the maximum age (in transactions) that a table's pg_class.relfrozenxid field can attain before a VACUUM operation is forced to prevent transaction ID wraparound within the table. Note that the system will launch autovacuum processes to prevent wraparound even when autovacuum is otherwise disabled.

Vacuum also allows removal of old files from the pg_xact subdirectory, which is why the default is a relatively low 200 million transactions. This parameter can only be set at server start, but the setting can be reduced for individual tables by changing table storage parameters. For more information see Section 24.1.5.

autovacuum_multixact_freeze_max_age(integer)

Specifies the maximum age (in multixacts) that a table's pg_class.relminmxid field can attain before a VACUUM operation is forced to prevent multixact ID wraparound within the table. Note that the system will launch autovacuum processes to prevent wraparound even when autovacuum is otherwise disabled.

Vacuuming multixacts also allows removal of old files from the pg_multixact/members and pg_multixact/offsets subdirectories, which is why the default is a relatively low 400 million multixacts. This parameter can only be set at server start, but the setting can be reduced for individual tables by changing table storage parameters. For more information see Section 24.1.5.1.

autovacuum_vacuum_cost_delay(integer)

Specifies the cost delay value that will be used in automatic VACUUM operations. If -1 is specified, the regular vacuum_cost_delay value will be used. The default value is 20 milliseconds. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_vacuum_cost_limit (integer)

Specifies the cost limit value that will be used in automatic VACUUM operations. If -1 is specified (which is the default), the regular vacuum_cost_limit value will be used. Note that the value is distributed proportionally among the running autovacuum workers, if there is more than one, so that the sum of the limits for each worker does not exceed the value of this variable. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

19.11. Client Connection Defaults

19.11.1. Statement Behavior

search_path (string)

This variable specifies the order in which schemas are searched when an object (table, data type, function, etc.) is referenced by a simple name with no schema specified. When there are objects of
identical names in different schemas, the one found first in the search path is used. An object that
is not in any of the schemas in the search path can only be referenced by specifying its containing
schema with a qualified (dotted) name.

The value for search_path must be a comma-separated list of schema names. Any name that
is not an existing schema, or is a schema for which the user does not have USAGE permission,
is silently ignored.

If one of the list items is the special name $user, then the schema having the name returned by
SESSION_USER is substituted, if there is such a schema and the user has USAGE permission
for it. (If not, $user is ignored.)

The system catalog schema, pg_catalog, is always searched, whether it is mentioned in the
path or not. If it is mentioned in the path then it will be searched in the specified order. If pg_cat-
alog is not in the path then it will be searched before searching any of the path items.

Likewise, the current session's temporary-table schema, pg_temp_nnn, is always searched if
it exists. It can be explicitly listed in the path by using the alias pg_temp. If it is not listed in
the path then it is searched first (even before pg_catalog). However, the temporary schema is
only searched for relation (table, view, sequence, etc) and data type names. It is never searched
for function or operator names.

When objects are created without specifying a particular target schema, they will be placed in the
first valid schema named in search_path. An error is reported if the search path is empty.

The default value for this parameter is "$user", public. This setting supports shared use of
a database (where no users have private schemas, and all share use of public), private per-user
schemas, and combinations of these. Other effects can be obtained by altering the default search
path setting, either globally or per-user.

For more information on schema handling, see Section 5.8. In particular, the default configuration
is suitable only when the database has a single user or a few mutually-trusting users.

The current effective value of the search path can be examined via the SQL function current_schemas
(see Section 9.25). This is not quite the same as examining the value of
search_path, since current_schemas shows how the items appearing in search_path
were resolved.

row_security (boolean)

This variable controls whether to raise an error in lieu of applying a row security policy. When
set to on, policies apply normally. When set to off, queries fail which would otherwise apply
at least one policy. The default is on. Change to off where limited row visibility could cause
incorrect results; for example, pg_dump makes that change by default. This variable has no ef-
fect on roles which bypass every row security policy, to wit, superusers and roles with the BY-
PASSRLS attribute.

For more information on row security policies, see CREATE POLICY.

default_tablespace (string)

This variable specifies the default tablespace in which to create objects (tables and indexes) when
a CREATE command does not explicitly specify a tablespace.

The value is either the name of a tablespace, or an empty string to specify using the default table-
space of the current database. If the value does not match the name of any existing tablespace,
PostgreSQL will automatically use the default tablespace of the current database. If a nondefault
tablespace is specified, the user must have CREATE privilege for it, or creation attempts will fail.

This variable is not used for temporary tables; for them, temp_tablespaces is consulted instead.
This variable is also not used when creating databases. By default, a new database inherits its tablespace setting from the template database it is copied from.

For more information on tablespaces, see Section 22.6.

temp_tablespaces (string)

This variable specifies tablespaces in which to create temporary objects (temp tables and indexes on temp tables) when a CREATE command does not explicitly specify a tablespace. Temporary files for purposes such as sorting large data sets are also created in these tablespaces.

The value is a list of names of tablespaces. When there is more than one name in the list, PostgreSQL chooses a random member of the list each time a temporary object is to be created; except that within a transaction, successively created temporary objects are placed in successive tablespaces from the list. If the selected element of the list is an empty string, PostgreSQL will automatically use the default tablespace of the current database instead.

When temp_tablespaces is set interactively, specifying a nonexistent tablespace is an error, as is specifying a tablespace for which the user does not have CREATE privilege. However, when using a previously set value, nonexistent tablespaces are ignored, as are tablespaces for which the user lacks CREATE privilege. In particular, this rule applies when using a value set in postgres.conf.

The default value is an empty string, which results in all temporary objects being created in the default tablespace of the current database.

See also default_tablespace.

check_function_bodies (boolean)

This parameter is normally on. When set to off, it disables validation of the function body string during CREATE FUNCTION. Disabling validation avoids side effects of the validation process and avoids false positives due to problems such as forward references. Set this parameter to off before loading functions on behalf of other users; pg_dump does so automatically.

default_transaction_isolation (enum)

Each SQL transaction has an isolation level, which can be either “read uncommitted”, “read committed”, “repeatable read”, or “serializable”. This parameter controls the default isolation level of each new transaction. The default is “read committed”.

Consult Chapter 13 and SET TRANSACTION for more information.

default_transaction_read_only (boolean)

A read-only SQL transaction cannot alter non-temporary tables. This parameter controls the default read-only status of each new transaction. The default is off (read/write).

Consult SET TRANSACTION for more information.

default_transaction_deferrable (boolean)

When running at the serializable isolation level, a deferrable read-only SQL transaction may be delayed before it is allowed to proceed. However, once it begins executing it does not incur any of the overhead required to ensure serializability; so serialization code will have no reason to force it to abort because of concurrent updates, making this option suitable for long-running read-only transactions.

This parameter controls the default deferrable status of each new transaction. It currently has no effect on read-write transactions or those operating at isolation levels lower than serializable. The default is off.
Server Configuration

Consult SET TRANSACTION for more information.

session_replication_role (enum)

Controls firing of replication-related triggers and rules for the current session. Setting this variable requires superuser privilege and results in discarding any previously cached query plans. Possible values are origin (the default), replica and local.

The intended use of this setting is that logical replication systems set it to replica when they are applying replicated changes. The effect of that will be that triggers and rules (that have not been altered from their default configuration) will not fire on the replica. See the ALTER TABLE clauses ENABLE TRIGGER and ENABLE RULE for more information.

PostgreSQL treats the settings origin and local the same internally. Third-party replication systems may use these two values for their internal purposes, for example using local to designate a session whose changes should not be replicated.

Since foreign keys are implemented as triggers, setting this parameter to replica also disables all foreign key checks, which can leave data in an inconsistent state if improperly used.

statement_timeout (integer)

Abort any statement that takes more than the specified number of milliseconds, starting from the time the command arrives at the server from the client. If log_min_error_statement is set to ERROR or lower, the statement that timed out will also be logged. A value of zero (the default) turns this off.

Setting statement_timeout in postgresql.conf is not recommended because it would affect all sessions.

lock_timeout (integer)

Abort any statement that waits longer than the specified number of milliseconds while attempting to acquire a lock on a table, index, row, or other database object. The time limit applies separately to each lock acquisition attempt. The limit applies both to explicit locking requests (such as LOCK TABLE, or SELECT FOR UPDATE without NOWAIT) and to implicitly-acquired locks. If log_min_error_statement is set to ERROR or lower, the statement that timed out will be logged. A value of zero (the default) turns this off.

Unlike statement_timeout, this timeout can only occur while waiting for locks. Note that if statement_timeout is nonzero, it is rather pointless to set lock_timeout to the same or larger value, since the statement timeout would always trigger first.

Setting lock_timeout in postgresql.conf is not recommended because it would affect all sessions.

idle_in_transaction_session_timeout (integer)

Terminate any session with an open transaction that has been idle for longer than the specified duration in milliseconds. This allows any locks held by that session to be released and the connection slot to be reused; it also allows tuples visible only to this transaction to be vacuumed. See Section 24.1 for more details about this.

The default value of 0 disables this feature.

vacuum_freeze_table_age (integer)

VACUUM performs an aggressive scan if the table's pg_class.relfrozenxid field has reached the age specified by this setting. An aggressive scan differs from a regular VACUUM in that it visits every page that might contain unfrozen XIDs or MXIDs, not just those that might contain dead tuples. The default is 150 million transactions. Although users can set this value
anywhere from zero to two billions, VACUUM will silently limit the effective value to 95% of autovacuum_freeze_max_age, so that a periodical manual VACUUM has a chance to run before an anti-wraparound autovacuum is launched for the table. For more information see Section 24.1.5.

vacuum_freeze_min_age (integer)

Specifies the cutoff age (in transactions) that VACUUM should use to decide whether to freeze row versions while scanning a table. The default is 50 million transactions. Although users can set this value anywhere from zero to one billion, VACUUM will silently limit the effective value to half the value of autovacuum_freeze_max_age, so that there is not an unreasonably short time between forced autovacuums. For more information see Section 24.1.5.

vacuum_multixact_freeze_table_age (integer)

VACUUM performs an aggressive scan if the table's pg_class.relminmxid field has reached the age specified by this setting. An aggressive scan differs from a regular VACUUM in that it visits every page that might contain unfrozen XIDs or MXIDs, not just those that might contain dead tuples. The default is 150 million multixacts. Although users can set this value anywhere from zero to two billions, VACUUM will silently limit the effective value to 95% of autovacuum_multixact_freeze_max_age, so that a periodical manual VACUUM has a chance to run before an anti-wraparound is launched for the table. For more information see Section 24.1.5.1.

vacuum_multixact_freeze_min_age (integer)

Specifies the cutoff age (in multixacts) that VACUUM should use to decide whether to replace multixact IDs with a newer transaction ID or multixact ID while scanning a table. The default is 5 million multixacts. Although users can set this value anywhere from zero to one billion, VACUUM will silently limit the effective value to half the value of autovacuum_multixact_freeze_max_age, so that there is not an unreasonably short time between forced autovacuums. For more information see Section 24.1.5.1.

vacuum_cleanup_index_scale_factor (floating point)

Specifies the fraction of the total number of heap tuples counted in the previous statistics collection that can be inserted without incurring an index scan at the VACUUM cleanup stage. This setting currently applies to B-tree indexes only.

If no tuples were deleted from the heap, B-tree indexes are still scanned at the VACUUM cleanup stage when at least one of the following conditions is met: the index statistics are stale, or the index contains deleted pages that can be recycled during cleanup. Index statistics are considered to be stale if the number of newly inserted tuples exceeds the vacuum_cleanup_index_scale_factor fraction of the total number of heap tuples detected by the previous statistics collection. The total number of heap tuples is stored in the index meta-page. Note that the meta-page does not include this data until VACUUM finds no dead tuples, so B-tree index scan at the cleanup stage can only be skipped if the second and subsequent VACUUM cycles detect no dead tuples.

The value can range from 0 to 10000000000. When vacuum_cleanup_index_scale_factor is set to 0, index scans are never skipped during VACUUM cleanup. The default value is 0.1.

bytea_output (enum)

Sets the output format for values of type bytea. Valid values are hex (the default) and escape (the traditional PostgreSQL format). See Section 8.4 for more information. The bytea type always accepts both formats on input, regardless of this setting.

xmlbinary (enum)

Sets how binary values are to be encoded in XML. This applies for example when bytea values are converted to XML by the functions xmlelement or xmlforest. Possible values are
base64 and hex, which are both defined in the XML Schema standard. The default is base64. For further information about XML-related functions, see Section 9.14.

The actual choice here is mostly a matter of taste, constrained only by possible restrictions in client applications. Both methods support all possible values, although the hex encoding will be somewhat larger than the base64 encoding.

\texttt{xml\_option(enum)}

Sets whether DOCUMENT or CONTENT is implicit when converting between XML and character string values. See Section 8.13 for a description of this. Valid values are DOCUMENT and CONTENT. The default is CONTENT.

According to the SQL standard, the command to set this option is

\texttt{SET XML OPTION \{ DOCUMENT | CONTENT \};}

This syntax is also available in PostgreSQL.

\texttt{gin\_pending\_list\_limit(integer)}

Sets the maximum size of the GIN pending list which is used when \texttt{fastupdate} is enabled. If the list grows larger than this maximum size, it is cleaned up by moving the entries in it to the main GIN data structure in bulk. The default is four megabytes (4MB). This setting can be overridden for individual GIN indexes by changing index storage parameters. See Section 66.4.1 and Section 66.5 for more information.

19.11.2. Locale and Formatting

\texttt{DateStyle(string)}

Sets the display format for date and time values, as well as the rules for interpreting ambiguous date input values. For historical reasons, this variable contains two independent components: the output format specification (ISO, Postgres, SQL, or German) and the input/output specification for year/month/day ordering (DMY, MDY, or YMD). These can be set separately or together. The keywords Euro and European are synonyms for DMY; the keywords US, NonEuro, and NonEuropean are synonyms for MDY. See Section 8.5 for more information. The built-in default is ISO, MDY, but \texttt{initdb} will initialize the configuration file with a setting that corresponds to the behavior of the chosen \texttt{lc\_time} locale.

\texttt{IntervalStyle(enum)}

Sets the display format for interval values. The value sql\_standard will produce output matching SQL standard interval literals. The value postgres (which is the default) will produce output matching PostgreSQL releases prior to 8.4 when the DateStyle parameter was set to ISO. The value postgres\_verbose will produce output matching PostgreSQL releases prior to 8.4 when the DateStyle parameter was set to non-ISO output. The value iso\_8601 will produce output matching the time interval “format with designators” defined in section 4.4.3.2 of ISO 8601.

The IntervalStyle parameter also affects the interpretation of ambiguous interval input. See Section 8.5.4 for more information.

\texttt{TimeZone(string)}

Sets the time zone for displaying and interpreting time stamps. The built-in default is GMT, but that is typically overridden in \texttt{postgresql.conf}; \texttt{initdb} will install a setting there corresponding to its system environment. See Section 8.5.3 for more information.
timezone_abbreviations (string)

Sets the collection of time zone abbreviations that will be accepted by the server for datetime input. The default is 'Default', which is a collection that works in most of the world; there are also 'Australia' and 'India', and other collections can be defined for a particular installation. See Section B.3 for more information.

extra_float_digits (integer)

This parameter adjusts the number of digits displayed for floating-point values, including float4, float8, and geometric data types. The parameter value is added to the standard number of digits (FLT_DIG or DBL_DIG as appropriate). The value can be set as high as 3, to include partially-significant digits; this is especially useful for dumping float data that needs to be restored exactly. Or it can be set negative to suppress unwanted digits. See also Section 8.1.3.

client_encoding (string)

Sets the client-side encoding (character set). The default is to use the database encoding. The character sets supported by the PostgreSQL server are described in Section 23.3.1.

lc_messages (string)

Sets the language in which messages are displayed. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

On some systems, this locale category does not exist. Setting this variable will still work, but there will be no effect. Also, there is a chance that no translated messages for the desired language exist. In that case you will continue to see the English messages.

Only superusers can change this setting, because it affects the messages sent to the server log as well as to the client, and an improper value might obscure the readability of the server logs.

lc_monetary (string)

Sets the locale to use for formatting monetary amounts, for example with the to_char family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

lc_numeric (string)

Sets the locale to use for formatting numbers, for example with the to_char family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

lc_time (string)

Sets the locale to use for formatting dates and times, for example with the to_char family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

default_text_search_config (string)

Selects the text search configuration that is used by those variants of the text search functions that do not have an explicit argument specifying the configuration. See Chapter 12 for further information. The built-in default is pg_catalog.simple, but initdb will initialize the config-
uration file with a setting that corresponds to the chosen lc_cctype locale, if a configuration matching that locale can be identified.

19.11.3. Shared Library Preloading

Several settings are available for preloading shared libraries into the server, in order to load additional functionality or achieve performance benefits. For example, a setting of '$libdir/mylib' would cause mylib.so (or on some platforms, mylib.sl) to be preloaded from the installation's standard library directory. The differences between the settings are when they take effect and what privileges are required to change them.

PostgreSQL procedural language libraries can be preloaded in this way, typically by using the syntax '$libdir/plXXX' where XXX is pgsql, perl, tcl, or python.

Only shared libraries specifically intended to be used with PostgreSQL can be loaded this way. Every PostgreSQL-supported library has a “magic block” that is checked to guarantee compatibility. For this reason, non-PostgreSQL libraries cannot be loaded in this way. You might be able to use operating-system facilities such as LD_PRELOAD for that.

In general, refer to the documentation of a specific module for the recommended way to load that module.

local_preload_libraries (string)

This variable specifies one or more shared libraries that are to be preloaded at connection start. It contains a comma-separated list of library names, where each name is interpreted as for the LOAD command. Whitespace between entries is ignored; surround a library name with double quotes if you need to include whitespace or commas in the name. The parameter value only takes effect at the start of the connection. Subsequent changes have no effect. If a specified library is not found, the connection attempt will fail.

This option can be set by any user. Because of that, the libraries that can be loaded are restricted to those appearing in the plugins subdirectory of the installation's standard library directory. (It is the database administrator's responsibility to ensure that only “safe” libraries are installed there.) Entries in local_preload_libraries can specify this directory explicitly, for example $libdir/plugins/mylib, or just specify the library name — mylib would have the same effect as $libdir/plugins/mylib.

The intent of this feature is to allow unprivileged users to load debugging or performance-measurement libraries into specific sessions without requiring an explicit LOAD command. To that end, it would be typical to set this parameter using the PGOPTIONS environment variable on the client or by using ALTER ROLE SET. However, unless a module is specifically designed to be used in this way by non-superusers, this is usually not the right setting to use. Look at session_preload_libraries instead.

session_preload_libraries (string)

This variable specifies one or more shared libraries that are to be preloaded at connection start. It contains a comma-separated list of library names, where each name is interpreted as for the LOAD command. Whitespace between entries is ignored; surround a library name with double quotes if you need to include whitespace or commas in the name. The parameter value only takes effect at the start of the connection. Subsequent changes have no effect. If a specified library is not found, the connection attempt will fail. Only superusers can change this setting.

The intent of this feature is to allow debugging or performance-measurement libraries to be loaded into specific sessions without an explicit LOAD command being given. For example, auto_explain could be enabled for all sessions under a given user name by setting this parameter with ALTER ROLE SET. Also, this parameter can be changed without restarting the server (but changes only
take effect when a new session is started), so it is easier to add new modules this way, even if they should apply to all sessions.

Unlike shared_preload_libraries, there is no large performance advantage to loading a library at session start rather than when it is first used. There is some advantage, however, when connection pooling is used.

shared_preload_libraries (string)

This variable specifies one or more shared libraries to be preloaded at server start. It contains a comma-separated list of library names, where each name is interpreted as for the LOAD command. Whitespace between entries is ignored; surround a library name with double quotes if you need to include whitespace or commas in the name. This parameter can only be set at server start. If a specified library is not found, the server will fail to start.

Some libraries need to perform certain operations that can only take place at postmaster start, such as allocating shared memory, reserving light-weight locks, or starting background workers. Those libraries must be loaded at server start through this parameter. See the documentation of each library for details.

Other libraries can also be preloaded. By preloading a shared library, the library startup time is avoided when the library is first used. However, the time to start each new server process might increase slightly, even if that process never uses the library. So this parameter is recommended only for libraries that will be used in most sessions. Also, changing this parameter requires a server restart, so this is not the right setting to use for short-term debugging tasks, say. Use session_preload_libraries for that instead.

Note

On Windows hosts, preloading a library at server start will not reduce the time required to start each new server process; each server process will re-load all preload libraries. However, shared_preload_libraries is still useful on Windows hosts for libraries that need to perform operations at postmaster start time.

jit_provider (string)

This variable is the name of the JIT provider library to be used (see Section 32.4.2). The default is llvmjit. This parameter can only be set at server start.

If set to a non-existent library, JIT will not be available, but no error will be raised. This allows JIT support to be installed separately from the main PostgreSQL package.

19.11.4. Other Defaults

dynamic_library_path (string)

If a dynamically loadable module needs to be opened and the file name specified in the CREATE FUNCTION or LOAD command does not have a directory component (i.e., the name does not contain a slash), the system will search this path for the required file.

The value for dynamic_library_path must be a list of absolute directory paths separated by colons (or semi-colons on Windows). If a list element starts with the special string $libdir, the compiled-in PostgreSQL package library directory is substituted for $libdir; this is where the modules provided by the standard PostgreSQL distribution are installed. (Use pg_config --pkglibdir to find out the name of this directory.) For example:
Server Configuration

dynamic_library_path = '/usr/local/lib/postgresql:/home/my_project/lib:$libdir'

or, in a Windows environment:

dynamic_library_path = 'C:\tools\postgresql;H:\my_project\lib;$libdir'

The default value for this parameter is '$libdir'. If the value is set to an empty string, the automatic path search is turned off.

This parameter can be changed at run time by superusers, but a setting done that way will only persist until the end of the client connection, so this method should be reserved for development purposes. The recommended way to set this parameter is in the `postgresql.conf` configuration file.

gin_fuzzy_search_limit (integer)

Soft upper limit of the size of the set returned by GIN index scans. For more information see Section 66.5.

19.12. Lock Management

deadlock_timeout (integer)

This is the amount of time, in milliseconds, to wait on a lock before checking to see if there is a deadlock condition. The check for deadlock is relatively expensive, so the server doesn't run it every time it waits for a lock. We optimistically assume that deadlocks are not common in production applications and just wait on the lock for a while before checking for a deadlock. Increasing this value reduces the amount of time wasted in needless deadlock checks, but slows down reporting of real deadlock errors. The default is one second (1s), which is probably about the smallest value you would want in practice. On a heavily loaded server you might want to raise it. Ideally the setting should exceed your typical transaction time, so as to improve the odds that a lock will be released before the waiter decides to check for deadlock. Only superusers can change this setting.

When log_lock_waits is set, this parameter also determines the length of time to wait before a log message is issued about the lock wait. If you are trying to investigate locking delays you might want to set a shorter than normal deadlock_timeout.

max_locks_per_transaction (integer)

The shared lock table tracks locks on max_locks_per_transaction * (max_connections + max_prepared_transactions) objects (e.g., tables); hence, no more than this many distinct objects can be locked at any one time. This parameter controls the average number of object locks allocated for each transaction; individual transactions can lock more objects as long as the locks of all transactions fit in the lock table. This is not the number of rows that can be locked; that value is unlimited. The default, 64, has historically proven sufficient, but you might need to raise this value if you have queries that touch many different tables in a single transaction, e.g. query of a parent table with many children. This parameter can only be set at server start.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

max_pred_locks_per_transaction (integer)

The shared predicate lock table tracks locks on max_pred_locks_per_transaction * (max_connections + max_prepared_transactions) objects (e.g., tables); hence, no more than this many distinct objects can be locked at any one time. This parameter controls the average number
Server Configuration

of object locks allocated for each transaction; individual transactions can lock more objects as long as the locks of all transactions fit in the lock table. This is not the number of rows that can be locked; that value is unlimited. The default, 64, has generally been sufficient in testing, but you might need to raise this value if you have clients that touch many different tables in a single serializable transaction. This parameter can only be set at server start.

max_pred_locks_per_relation (integer)

This controls how many pages or tuples of a single relation can be predicate-locked before the lock is promoted to covering the whole relation. Values greater than or equal to zero mean an absolute limit, while negative values mean max_pred_locks_per_transaction divided by the absolute value of this setting. The default is -2, which keeps the behavior from previous versions of PostgreSQL. This parameter can only be set in the postgresql.conf file or on the server command line.

max_pred_locks_per_page (integer)

This controls how many rows on a single page can be predicate-locked before the lock is promoted to covering the whole page. The default is 2. This parameter can only be set in the postgresql.conf file or on the server command line.

19.13. Version and Platform Compatibility

19.13.1. Previous PostgreSQL Versions

array_nulls (boolean)

This controls whether the array input parser recognizes unquoted NULL as specifying a null array element. By default, this is on, allowing array values containing null values to be entered. However, PostgreSQL versions before 8.2 did not support null values in arrays, and therefore would treat NULL as specifying a normal array element with the string value “NULL”. For backward compatibility with applications that require the old behavior, this variable can be turned off.

Note that it is possible to create array values containing null values even when this variable is off.

backslash_quote (enum)

This controls whether a quote mark can be represented by \' in a string literal. The preferred, SQL-standard way to represent a quote mark is by doubling it (""') but PostgreSQL has historically also accepted ". However, use of \' creates security risks because in some client character set encodings, there are multibyte characters in which the last byte is numerically equivalent to ASCII \. If client-side code does escaping incorrectly then a SQL-injection attack is possible. This risk can be prevented by making the server reject queries in which a quote mark appears to be escaped by a backslash. The allowed values of backslash_quote are on (allow \' always), off (reject always), and safe_encoding (allow only if client encoding does not allow ASCII \ within a multibyte character). safe_encoding is the default setting.

Note that in a standard-conforming string literal, \ just means \ anyway. This parameter only affects the handling of non-standard-conforming literals, including escape string syntax (E'...').

default_with_oids (boolean)

This controls whether CREATE TABLE and CREATE TABLE AS include an OID column in newly-created tables, if neither WITH OIDS nor WITHOUT OIDS is specified. It also determines whether OIDs will be included in tables created by SELECT INTO. The parameter is off by default; in PostgreSQL 8.0 and earlier, it was on by default.

The use of OIDs in user tables is considered deprecated, so most installations should leave this variable disabled. Applications that require OIDs for a particular table should specify WITH OIDS when creating the table. This variable can be enabled for compatibility with old applications that do not follow this behavior.
escape_string_warning (boolean)

When on, a warning is issued if a backslash (\) appears in an ordinary string literal ("...") syntax and standard_conforming_strings is off. The default is on.

Applications that wish to use backslash as escape should be modified to use escape string syntax (E"..."), because the default behavior of ordinary strings is now to treat backslash as an ordinary character, per SQL standard. This variable can be enabled to help locate code that needs to be changed.

lo_compat_privileges (boolean)

In PostgreSQL releases prior to 9.0, large objects did not have access privileges and were, therefore, always readable and writable by all users. Setting this variable to on disables the new privilege checks, for compatibility with prior releases. The default is off. Only superusers can change this setting.

Setting this variable does not disable all security checks related to large objects — only those for which the default behavior has changed in PostgreSQL 9.0.

operator_precedence_warning (boolean)

When on, the parser will emit a warning for any construct that might have changed meanings since PostgreSQL 9.4 as a result of changes in operator precedence. This is useful for auditing applications to see if precedence changes have broken anything; but it is not meant to be kept turned on in production, since it will warn about some perfectly valid, standard-compliant SQL code. The default is off.

See Section 4.1.6 for more information.

quote_all_identifiers (boolean)

When the database generates SQL, force all identifiers to be quoted, even if they are not (currently) keywords. This will affect the output of EXPLAIN as well as the results of functions like pg_get_viewdef. See also the --quote-all-identifiers option of pg_dump and pg_dumpall.

standard_conforming_strings (boolean)

This controls whether ordinary string literals ("...") treat backslashes literally, as specified in the SQL standard. Beginning in PostgreSQL 9.1, the default is on (prior releases defaulted to off). Applications can check this parameter to determine how string literals will be processed. The presence of this parameter can also be taken as an indication that the escape string syntax (E"...") is supported. Escape string syntax (Section 4.1.2.2) should be used if an application desires backslashes to be treated as escape characters.

synchronize_seqscans (boolean)

This allows sequential scans of large tables to synchronize with each other, so that concurrent scans read the same block at about the same time and hence share the I/O workload. When this is enabled, a scan might start in the middle of the table and then "wrap around" the end to cover all rows, so as to synchronize with the activity of scans already in progress. This can result in unpredictable changes in the row ordering returned by queries that have no ORDER BY clause. Setting this parameter to off ensures the pre-8.3 behavior in which a sequential scan always starts from the beginning of the table. The default is on.

19.13.2. Platform and Client Compatibility

transform_null_equals (boolean)

When on, expressions of the form expr = NULL (or NULL = expr) are treated as expr IS NULL, that is, they return true if expr evaluates to the null value, and false otherwise. The
correct SQL-spec-compliant behavior of $expr = \text{NULL}$ is to always return null (unknown). Therefore this parameter defaults to off.

However, filtered forms in Microsoft Access generate queries that appear to use $expr = \text{NULL}$ to test for null values, so if you use that interface to access the database you might want to turn this option on. Since expressions of the form $expr = \text{NULL}$ always return the null value (using the SQL standard interpretation), they are not very useful and do not appear often in normal applications so this option does little harm in practice. But new users are frequently confused about the semantics of expressions involving null values, so this option is off by default.

Note that this option only affects the exact form $= \text{NULL}$, not other comparison operators or other expressions that are computationally equivalent to some expression involving the equals operator (such as $\text{IN}$). Thus, this option is not a general fix for bad programming.

Refer to Section 9.2 for related information.

19.14. Error Handling

exit_on_error(boolean)

If true, any error will terminate the current session. By default, this is set to false, so that only FATAL errors will terminate the session.

restart_after_crash(boolean)

When set to true, which is the default, PostgreSQL will automatically reinitialize after a backend crash. Leaving this value set to true is normally the best way to maximize the availability of the database. However, in some circumstances, such as when PostgreSQL is being invoked by clusterware, it may be useful to disable the restart so that the clusterware can gain control and take any actions it deems appropriate.

19.15. Preset Options

The following “parameters” are read-only, and are determined when PostgreSQL is compiled or when it is installed. As such, they have been excluded from the sample postgresql.conf file. These options report various aspects of PostgreSQL behavior that might be of interest to certain applications, particularly administrative front-ends.

block_size(integer)

Reports the size of a disk block. It is determined by the value of BLCKSZ when building the server. The default value is 8192 bytes. The meaning of some configuration variables (such as shared_buffers) is influenced by block_size. See Section 19.4 for information.

data_checksums(boolean)

Reports whether data checksums are enabled for this cluster. See data checksums for more information.

data_directory_mode(integer)

On Unix systems this parameter reports the permissions of the data directory defined by (data_directory) at startup. (On Microsoft Windows this parameter will always display 0700). See group access for more information.

debug_assertions(boolean)

Reports whether PostgreSQL has been built with assertions enabled. That is the case if the macro USE_ASSERT_CHECKING is defined when PostgreSQL is built (accomplished e.g. by the configure option --enable-cassert). By default PostgreSQL is built without assertions.
integer_datetimes (boolean)
Reports whether PostgreSQL was built with support for 64-bit-integer dates and times. As of PostgreSQL 10, this is always on.

lc_collate (string)
Reports the locale in which sorting of textual data is done. See Section 23.1 for more information. This value is determined when a database is created.

lc_cctype (string)
Reports the locale that determines character classifications. See Section 23.1 for more information. This value is determined when a database is created. Ordinarily this will be the same as lc_collate, but for special applications it might be set differently.

max_function_args (integer)
Reports the maximum number of function arguments. It is determined by the value of FUNC_MAX_ARGS when building the server. The default value is 100 arguments.

max_identifier_length (integer)
Reports the maximum identifier length. It is determined as one less than the value of NAME_DATALEN when building the server. The default value of NAME_DATALEN is 64; therefore the default max_identifier_length is 63 bytes, which can be less than 63 characters when using multibyte encodings.

max_index_keys (integer)
Reports the maximum number of index keys. It is determined by the value of INDEX_MAX_KEYS when building the server. The default value is 32 keys.

segment_size (integer)
Reports the number of blocks (pages) that can be stored within a file segment. It is determined by the value of RELSEG_SIZE when building the server. The maximum size of a segment file in bytes is equal to segment_size multiplied by block_size; by default this is 1GB.

server_encoding (string)
Reports the database encoding (character set). It is determined when the database is created. Ordinarily, clients need only be concerned with the value of client_encoding.

server_version (string)
Reports the version number of the server. It is determined by the value of PG_VERSION when building the server.

server_version_num (integer)
Reports the version number of the server as an integer. It is determined by the value of PG_VERSION_NUM when building the server.

wal_block_size (integer)
Reports the size of a WAL disk block. It is determined by the value of XLOG_BLCKSZ when building the server. The default value is 8192 bytes.

wal_segment_size (integer)
Reports the number of blocks (pages) in a WAL segment file. The total size of a WAL segment file in bytes is equal to wal_segment_size multiplied by wal_block_size; by default this is 16MB. See Section 30.4 for more information.
19.16. Customized Options

This feature was designed to allow parameters not normally known to PostgreSQL to be added by add-on modules (such as procedural languages). This allows extension modules to be configured in the standard ways.

Custom options have two-part names: an extension name, then a dot, then the parameter name proper, much like qualified names in SQL. An example is `plpgsql.variable_conflict`.

Because custom options may need to be set in processes that have not loaded the relevant extension module, PostgreSQL will accept a setting for any two-part parameter name. Such variables are treated as placeholders and have no function until the module that defines them is loaded. When an extension module is loaded, it will add its variable definitions, convert any placeholder values according to those definitions, and issue warnings for any unrecognized placeholders that begin with its extension name.

19.17. Developer Options

The following parameters are intended for work on the PostgreSQL source code, and in some cases to assist with recovery of severely damaged databases. There should be no reason to use them on a production database. As such, they have been excluded from the sample `postgresql.conf` file. Note that many of these parameters require special source compilation flags to work at all.

`allow_system_table_mods(boolean)`

Allows modification of the structure of system tables. This is used by `initdb`. This parameter can only be set at server start.

`ignore_system_indexes(boolean)`

Ignore system indexes when reading system tables (but still update the indexes when modifying the tables). This is useful when recovering from damaged system indexes. This parameter cannot be changed after session start.

`post_auth_delay(integer)`

If nonzero, a delay of this many seconds occurs when a new server process is started, after it conducts the authentication procedure. This is intended to give developers an opportunity to attach to the server process with a debugger. This parameter cannot be changed after session start.

`pre_auth_delay(integer)`

If nonzero, a delay of this many seconds occurs just after a new server process is forked, before it conducts the authentication procedure. This is intended to give developers an opportunity to attach to the server process with a debugger to trace down misbehavior in authentication. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`trace_notify(boolean)`

Generates a great amount of debugging output for the `LISTEN` and `NOTIFY` commands. `client_min_messages` or `log_min_messages` must be `DEBUG1` or lower to send this output to the client or server logs, respectively.

`trace_recovery_messages(enum)`

Enables logging of recovery-related debugging output that otherwise would not be logged. This parameter allows the user to override the normal setting of `log_min_messages`, but only for specific messages. This is intended for use in debugging Hot Standby. Valid values are `DEBUG5`, `DEBUG4`, `DEBUG3`, `DEBUG2`, `DEBUG1`, and `LOG`. The default, `LOG`, does not affect logging decisions at all. The other values cause recovery-related debug messages of that priority or higher.
to be logged as though they had LOG priority; for common settings of log_min_messages this results in unconditionally sending them to the server log. This parameter can only be set in the postgresql.conf file or on the server command line.

**trace_sort** (boolean)

If on, emit information about resource usage during sort operations. This parameter is only available if the TRACE_SORT macro was defined when PostgreSQL was compiled. (However, TRACE_SORT is currently defined by default.)

**trace_locks** (boolean)

If on, emit information about lock usage. Information dumped includes the type of lock operation, the type of lock and the unique identifier of the object being locked or unlocked. Also included are bit masks for the lock types already granted on this object as well as for the lock types awaited on this object. For each lock type a count of the number of granted locks and waiting locks is also dumped as well as the totals. An example of the log file output is shown here:

```
LOG:  LockAcquire: new: lock(0xb7acd844)  id(24688,24696,0,0,0,1)
  grantMask(0)  req(0,0,0,0,0,0)=0  grant(0,0,0,0,0,0)=0
  wait(0)  type(AccessShareLock)
LOG:  GrantLock: lock(0xb7acd844)  id(24688,24696,0,0,0,1)
  grantMask(2)  req(1,0,0,0,0,0)=1  grant(1,0,0,0,0,0)=1
  wait(0)  type(AccessShareLock)
LOG:  UnGrantLock: updated: lock(0xb7acd844)
  id(24688,24696,0,0,0,1)
  grantMask(0)  req(0,0,0,0,0,0)=0  grant(0,0,0,0,0,0)=0
  wait(0)  type(AccessShareLock)
LOG:  CleanUpLock: deleting: lock(0xb7acd844)
  id(24688,24696,0,0,0,1)
  grantMask(0)  req(0,0,0,0,0,0)=0  grant(0,0,0,0,0,0)=0
  wait(0)  type(INVALID)
```

Details of the structure being dumped may be found in src/include/storage/lock.h.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

**trace_lwlocks** (boolean)

If on, emit information about lightweight lock usage. Lightweight locks are intended primarily to provide mutual exclusion of access to shared-memory data structures.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

**trace_userlocks** (boolean)

If on, emit information about user lock usage. Output is the same as for trace_locks, only for advisory locks.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

**trace_lock_oidmin** (integer)

If set, do not trace locks for tables below this OID. (use to avoid output on system tables)

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.
Server Configuration

trace_lock_table(integer)

Unconditionally trace locks on this table (OID).

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

dump_deadlocks(boolean)

If set, dumps information about all current locks when a deadlock timeout occurs.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

log_btree_build_stats(boolean)

If set, logs system resource usage statistics (memory and CPU) on various B-tree operations.

This parameter is only available if the BTREE_BUILD_STATS macro was defined when PostgreSQL was compiled.

wal_consistency_checking(string)

This parameter is intended to be used to check for bugs in the WAL redo routines. When enabled, full-page images of any buffers modified in conjunction with the WAL record are added to the record. If the record is subsequently replayed, the system will first apply each record and then test whether the buffers modified by the record match the stored images. In certain cases (such as hint bits), minor variations are acceptable, and will be ignored. Any unexpected differences will result in a fatal error, terminating recovery.

The default value of this setting is the empty string, which disables the feature. It can be set to all to check all records, or to a comma-separated list of resource managers to check only records originating from those resource managers. Currently, the supported resource managers are heap, heap2, btree, hash, gin, gist, sequence, spgist, brin, and generic. Only superusers can change this setting.

wal_debug(boolean)

If on, emit WAL-related debugging output. This parameter is only available if the WAL_DEBUG macro was defined when PostgreSQL was compiled.

ignore_checksum_failure(boolean)

Only has effect if data checksums are enabled.

Detection of a checksum failure during a read normally causes PostgreSQL to report an error, aborting the current transaction. Setting ignore_checksum_failure to on causes the system to ignore the failure (but still report a warning), and continue processing. This behavior may cause crashes, propagate or hide corruption, or other serious problems. However, it may allow you to get past the error and retrieve undamaged tuples that might still be present in the table if the block header is still sane. If the header is corrupt an error will be reported even if this option is enabled. The default setting is off, and it can only be changed by a superuser.

zero_damaged_pages(boolean)

Detection of a damaged page header normally causes PostgreSQL to report an error, aborting the current transaction. Setting zero_damaged_pages to on causes the system to instead report a warning, zero out the damaged page in memory, and continue processing. This behavior will destroy data, namely all the rows on the damaged page. However, it does allow you to get past the error and retrieve rows from any undamaged pages that might be present in the table. It is useful for recovering data if corruption has occurred due to a hardware or software error. You should
generally not set this on until you have given up hope of recovering data from the damaged pages of a table. Zeroed-out pages are not forced to disk so it is recommended to recreate the table or the index before turning this parameter off again. The default setting is off, and it can only be changed by a superuser.

**jit_debugging_support (boolean)**

If LLVM has the required functionality, register generated functions with GDB. This makes debugging easier. The default setting is off. This parameter can only be set at server start.

**jit_dump_bitcode (boolean)**

Writes the generated LLVM IR out to the file system, inside data_directory. This is only useful for working on the internals of the JIT implementation. The default setting is off. This parameter can only be changed by a superuser.

**jit_expressions (boolean)**

Determines whether expressions are JIT compiled, when JIT compilation is activated (see Section 32.2). The default is on.

**jit_profiling_support (boolean)**

If LLVM has the required functionality, emit the data needed to allow perf to profile functions generated by JIT. This writes out files to $HOME/.debug/jit/; the user is responsible for performing cleanup when desired. The default setting is off. This parameter can only be set at server start.

**jit_tuple_deforming (boolean)**

Determines whether tuple deforming is JIT compiled, when JIT compilation is activated (see Section 32.2). The default is on.

## 19.18. Short Options

For convenience there are also single letter command-line option switches available for some parameters. They are described in Table 19.2. Some of these options exist for historical reasons, and their presence as a single-letter option does not necessarily indicate an endorsement to use the option heavily.

### Table 19.2. Short Option Key

<table>
<thead>
<tr>
<th>Short Option</th>
<th>Equivalent</th>
</tr>
</thead>
<tbody>
<tr>
<td>-B x</td>
<td>shared_buffers = x</td>
</tr>
<tr>
<td>-d x</td>
<td>log_min_messages = DEBUGx</td>
</tr>
<tr>
<td>-e</td>
<td>datestyle = euro</td>
</tr>
<tr>
<td>-fb, -fh, -fi, -fm, -fn, -fo, -fs, -ft</td>
<td>enable_bitmapscan = off, enable_hashjoin = off, enable_indexscan = off, enable_mergejoin = off, enable_nestloop = off, enable_indexonlyscan = off, enable_seqscan = off, enable_tidscan = off</td>
</tr>
<tr>
<td>-F</td>
<td>fsync = off</td>
</tr>
<tr>
<td>-h x</td>
<td>listen_addresses = x</td>
</tr>
<tr>
<td>-i</td>
<td>listen_addresses = '*'</td>
</tr>
<tr>
<td>Short Option</td>
<td>Equivalent</td>
</tr>
<tr>
<td>--------------</td>
<td>------------</td>
</tr>
<tr>
<td>-k x</td>
<td>unix_socket_directories = x</td>
</tr>
<tr>
<td>-l</td>
<td>ssl = on</td>
</tr>
<tr>
<td>-N x</td>
<td>max_connections = x</td>
</tr>
<tr>
<td>-O</td>
<td>allow_system_table_mods = on</td>
</tr>
<tr>
<td>-p x</td>
<td>port = x</td>
</tr>
<tr>
<td>-P</td>
<td>ignore_system_indexes = on</td>
</tr>
<tr>
<td>-s</td>
<td>log_statement_stats = on</td>
</tr>
<tr>
<td>-S x</td>
<td>work_mem = x</td>
</tr>
<tr>
<td>-tpa,-tpl,-te</td>
<td>log_parser_stats = on, log_planner_stats = on, log_executor_stats = on</td>
</tr>
<tr>
<td>-W x</td>
<td>post_auth_delay = x</td>
</tr>
</tbody>
</table>
Chapter 20. Client Authentication

When a client application connects to the database server, it specifies which PostgreSQL database user name it wants to connect as, much the same way one logs into a Unix computer as a particular user. Within the SQL environment the active database user name determines access privileges to database objects — see Chapter 21 for more information. Therefore, it is essential to restrict which database users can connect.

Note

As explained in Chapter 21, PostgreSQL actually does privilege management in terms of “roles”. In this chapter, we consistently use database user to mean “role with the LOGIN privilege”.

Authentication is the process by which the database server establishes the identity of the client, and by extension determines whether the client application (or the user who runs the client application) is permitted to connect with the database user name that was requested.

PostgreSQL offers a number of different client authentication methods. The method used to authenticate a particular client connection can be selected on the basis of (client) host address, database, and user.

PostgreSQL database user names are logically separate from user names of the operating system in which the server runs. If all the users of a particular server also have accounts on the server’s machine, it makes sense to assign database user names that match their operating system user names. However, a server that accepts remote connections might have many database users who have no local operating system account, and in such cases there need be no connection between database user names and OS user names.

20.1. The pg_hba.conf File

Client authentication is controlled by a configuration file, which traditionally is named pg_hba.conf and is stored in the database cluster’s data directory. (HBA stands for host-based authentication.) A default pg_hba.conf file is installed when the data directory is initialized by initdb. It is possible to place the authentication configuration file elsewhere, however; see the hba_file configuration parameter.

The general format of the pg_hba.conf file is a set of records, one per line. Blank lines are ignored, as is any text after the # comment character. Records cannot be continued across lines. A record is made up of a number of fields which are separated by spaces and/or tabs. Fields can contain white space if the field value is double-quoted. Quoting one of the keywords in a database, user, or address field (e.g., all or replication) makes the word lose its special meaning, and just match a database, user, or host with that name.

Each record specifies a connection type, a client IP address range (if relevant for the connection type), a database name, a user name, and the authentication method to be used for connections matching these parameters. The first record with a matching connection type, client address, requested database, and user name is used to perform authentication. There is no “fall-through” or “backup”: if one record is chosen and the authentication fails, subsequent records are not considered. If no record matches, access is denied.

A record can have one of the seven formats

local  database user  auth-method [auth-options]
host   database user  address  auth-method [auth-options]
hostssl database user  address  auth-method [auth-options]
Client Authentication

hostnossl  database  user  address  auth-method  [auth-options]
host  database  user  IP-address  IP-mask  auth-method  [auth-options]
hostssl  database  user  IP-address  IP-mask  auth-method  [auth-options]
hostnossl  database  user  IP-address  IP-mask  auth-method  [auth-options]

The meaning of the fields is as follows:

local
This record matches connection attempts using Unix-domain sockets. Without a record of this type, Unix-domain socket connections are disallowed.

host
This record matches connection attempts made using TCP/IP. host records match either SSL or non-SSL connection attempts.

Note
Remote TCP/IP connections will not be possible unless the server is started with an appropriate value for the listen_addresses configuration parameter, since the default behavior is to listen for TCP/IP connections only on the local loopback address localhost.

hostssl
This record matches connection attempts made using TCP/IP, but only when the connection is made with SSL encryption.

To make use of this option the server must be built with SSL support. Furthermore, SSL must be enabled by setting the ssl configuration parameter (see Section 18.9 for more information). Otherwise, the hostssl record is ignored except for logging a warning that it cannot match any connections.

hostnossl
This record type has the opposite behavior of hostssl; it only matches connection attempts made over TCP/IP that do not use SSL.

database
Specifies which database name(s) this record matches. The value all specifies that it matches all databases. The value sameuser specifies that the record matches if the requested database has the same name as the requested user. The value samerole specifies that the requested user must be a member of the role with the same name as the requested database. (samegroup is an obsolete but still accepted spelling of samerole.) Superusers are not considered to be members of a role for the purposes of samerole unless they are explicitly members of the role, directly or indirectly, and not just by virtue of being a superuser. The value replication specifies that the record matches if a physical replication connection is requested (note that replication connections do not specify any particular database). Otherwise, this is the name of a specific PostgreSQL database. Multiple database names can be supplied by separating them with commas. A separate file containing database names can be specified by preceding the file name with @.

user
Specifies which database user name(s) this record matches. The value all specifies that it matches all users. Otherwise, this is either the name of a specific database user, or a group name preceded by +. (Recall that there is no real distinction between users and groups in PostgreSQL; a
Client Authentication

+ mark really means “match any of the roles that are directly or indirectly members of this role”, while a name without a + mark matches only that specific role.) For this purpose, a superuser is only considered to be a member of a role if they are explicitly a member of the role, directly or indirectly, and not just by virtue of being a superuser. Multiple user names can be supplied by separating them with commas. A separate file containing user names can be specified by preceding the file name with @.

address

Specifies the client machine address(es) that this record matches. This field can contain either a host name, an IP address range, or one of the special key words mentioned below.

An IP address range is specified using standard numeric notation for the range's starting address, then a slash (/) and a CIDR mask length. The mask length indicates the number of high-order bits of the client IP address that must match. Bits to the right of this should be zero in the given IP address. There must not be any white space between the IP address, the /, and the CIDR mask length.

Typical examples of an IPv4 address range specified this way are 172.20.143.89/32 for a single host, or 172.20.143.0/24 for a small network, or 10.6.0.0/16 for a larger one. An IPv6 address range might look like ::1/128 for a single host (in this case the IPv6 loopback address) or fe80::7a31:c1ff:0000:0000/96 for a small network. 0.0.0.0/0 represents all IPv4 addresses, and ::/0 represents all IPv6 addresses. To specify a single host, use a mask length of 32 for IPv4 or 128 for IPv6. In a network address, do not omit trailing zeroes.

An entry given in IPv4 format will match only IPv4 connections, and an entry given in IPv6 format will match only IPv6 connections, even if the represented address is in the IPv4-in-IPv6 range. Note that entries in IPv6 format will be rejected if the system's C library does not have support for IPv6 addresses.

You can also write all to match any IP address, samehost to match any of the server's own IP addresses, or samenet to match any address in any subnet that the server is directly connected to.

If a host name is specified (anything that is not an IP address range or a special key word is treated as a host name), that name is compared with the result of a reverse name resolution of the client's IP address (e.g., reverse DNS lookup, if DNS is used). Host name comparisons are case insensitive. If there is a match, then a forward name resolution (e.g., forward DNS lookup) is performed on the host name to check whether any of the addresses it resolves to are equal to the client's IP address. If both directions match, then the entry is considered to match. (The host name that is used in pg_hba.conf should be the one that address-to-name resolution of the client's IP address returns, otherwise the line won't be matched. Some host name databases allow associating an IP address with multiple host names, but the operating system will only return one host name when asked to resolve an IP address.)

A host name specification that starts with a dot (.) matches a suffix of the actual host name. So .example.com would match foo.example.com (but not just example.com).

When host names are specified in pg_hba.conf, you should make sure that name resolution is reasonably fast. It can be of advantage to set up a local name resolution cache such as nscd. Also, you may wish to enable the configuration parameter log_hostname to see the client's host name instead of the IP address in the log.

This field only applies to host, hostssl, and hostnossl records.

Note

Users sometimes wonder why host names are handled in this seemingly complicated way, with two name resolutions including a reverse lookup of the client's IP address. This complicates use of the feature in case the client's reverse DNS entry is not set up or yields some undesirable host name. It is done primarily for
efficiency: this way, a connection attempt requires at most two resolver lookups, one reverse and one forward. If there is a resolver problem with some address, it becomes only that client's problem. A hypothetical alternative implementation that only did forward lookups would have to resolve every host name mentioned in `pg_hba.conf` during every connection attempt. That could be quite slow if many names are listed. And if there is a resolver problem with one of the host names, it becomes everyone's problem.

Also, a reverse lookup is necessary to implement the suffix matching feature, because the actual client host name needs to be known in order to match it against the pattern.

Note that this behavior is consistent with other popular implementations of host name-based access control, such as the Apache HTTP Server and TCP Wrappers.

<table>
<thead>
<tr>
<th>IP-address</th>
<th>IP-mask</th>
</tr>
</thead>
<tbody>
<tr>
<td>These two fields can be used as an alternative to the <code>IP-address/mask-length</code> notation. Instead of specifying the mask length, the actual mask is specified in a separate column. For example, <code>255.0.0.0</code> represents an IPv4 CIDR mask length of 8, and <code>255.255.255.255</code> represents a CIDR mask length of 32.</td>
<td></td>
</tr>
<tr>
<td>These fields only apply to <code>host</code>, <code>hostssl</code>, and <code>hostnossl</code> records.</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>auth-method</th>
</tr>
</thead>
<tbody>
<tr>
<td>Specifies the authentication method to use when a connection matches this record. The possible choices are summarized here; details are in Section 20.3.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>trust</th>
</tr>
</thead>
<tbody>
<tr>
<td>Allow the connection unconditionally. This method allows anyone that can connect to the PostgreSQL database server to login as any PostgreSQL user they wish, without the need for a password or any other authentication. See Section 20.4 for details.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>reject</th>
</tr>
</thead>
<tbody>
<tr>
<td>Reject the connection unconditionally. This is useful for “filtering out” certain hosts from a group, for example a <code>reject</code> line could block a specific host from connecting, while a later line allows the remaining hosts in a specific network to connect.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>scram-sha-256</th>
</tr>
</thead>
<tbody>
<tr>
<td>Perform SCRAM-SHA-256 authentication to verify the user's password. See Section 20.5 for details.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>md5</th>
</tr>
</thead>
<tbody>
<tr>
<td>Perform SCRAM-SHA-256 or MD5 authentication to verify the user's password. See Section 20.5 for details.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>password</th>
</tr>
</thead>
<tbody>
<tr>
<td>Require the client to supply an unencrypted password for authentication. Since the password is sent in clear text over the network, this should not be used on untrusted networks. See Section 20.5 for details.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>gss</th>
</tr>
</thead>
<tbody>
<tr>
<td>Use GSSAPI to authenticate the user. This is only available for TCP/IP connections. See Section 20.6 for details.</td>
</tr>
</tbody>
</table>
Client Authentication

sspi

Use SSPI to authenticate the user. This is only available on Windows. See Section 20.7 for details.

ident

Obtain the operating system user name of the client by contacting the ident server on the client and check if it matches the requested database user name. Ident authentication can only be used on TCP/IP connections. When specified for local connections, peer authentication will be used instead. See Section 20.8 for details.

peer

Obtain the client's operating system user name from the operating system and check if it matches the requested database user name. This is only available for local connections. See Section 20.9 for details.

ldap

Authenticate using an LDAP server. See Section 20.10 for details.

radius

Authenticate using a RADIUS server. See Section 20.11 for details.

cert

Authenticate using SSL client certificates. See Section 20.12 for details.

pam

Authenticate using the Pluggable Authentication Modules (PAM) service provided by the operating system. See Section 20.13 for details.

bsd

Authenticate using the BSD Authentication service provided by the operating system. See Section 20.14 for details.

auth-options

After the auth-method field, there can be field(s) of the form name=value that specify options for the authentication method. Details about which options are available for which authentication methods appear below.

In addition to the method-specific options listed below, there is one method-independent authentication option clientcert, which can be specified in any hostssl record. When set to 1, this option requires the client to present a valid (trusted) SSL certificate, in addition to the other requirements of the authentication method.

Files included by @ constructs are read as lists of names, which can be separated by either whitespace or commas. Comments are introduced by #, just as in pg_hba.conf, and nested @ constructs are allowed. Unless the file name following @ is an absolute path, it is taken to be relative to the directory containing the referencing file.

Since the pg_hba.conf records are examined sequentially for each connection attempt, the order of the records is significant. Typically, earlier records will have tight connection match parameters and weaker authentication methods, while later records will have looser match parameters and stronger authentication methods. For example, one might wish to use trust authentication for local TCP/IP connections but require a password for remote TCP/IP connections. In this case a record specifying trust authentication for connections from 127.0.0.1 would appear before a record specifying password authentication for a wider range of allowed client IP addresses.
The pg_hba.conf file is read on start-up and when the main server process receives a SIGHUP signal. If you edit the file on an active system, you will need to signal the postmaster (using `pg_ctl` reload or `kill -HUP`) to make it re-read the file.

**Note**

The preceding statement is not true on Microsoft Windows: there, any changes in the pg_hba.conf file are immediately applied by subsequent new connections.

The system view `pg_hba_file_rules` can be helpful for pre-testing changes to the pg_hba.conf file, or for diagnosing problems if loading of the file did not have the desired effects. Rows in the view with non-null error fields indicate problems in the corresponding lines of the file.

**Tip**

To connect to a particular database, a user must not only pass the pg_hba.conf checks, but must have the CONNECT privilege for the database. If you wish to restrict which users can connect to which databases, it's usually easier to control this by granting/revoking CONNECT privilege than to put the rules in pg_hba.conf entries.

Some examples of pg_hba.conf entries are shown in Example 20.1. See the next section for details on the different authentication methods.

**Example 20.1. Example pg_hba.conf Entries**

```
# Allow any user on the local system to connect to any database
# with any database user name using Unix-domain sockets (the default for local
# connections).
#
# TYPE  DATABASE        USER            ADDRESS
METHOD
local   all             all             trust

# The same using local loopback TCP/IP connections.
#
# TYPE  DATABASE        USER            ADDRESS
METHOD
host    all             all             127.0.0.1/32
trust

# The same as the previous line, but using a separate netmask column
#
# TYPE  DATABASE        USER            IP-ADDRESS      IP-MASK
METHOD
host    all             all             127.0.0.1
255.255.255.255     trust

# The same over IPv6.
#
# TYPE  DATABASE        USER            ADDRESS
METHOD
```
Client Authentication

```plaintext

host    all             all             ::1/128
trust

# The same using a host name (would typically cover both IPv4 and IPv6).
# TYPE  DATABASE        USER            ADDRESS          METHOD
host    all             all             localhost
trust

# Allow any user from any host with IP address 192.168.93.x to connect
to database "postgres" as the same user name that ident reports for
the connection (typically the operating system user name).
# TYPE  DATABASE        USER            ADDRESS          METHOD
host    postgres        all             192.168.93.0/24
ident

# Allow any user from host 192.168.12.10 to connect to database
"postgres" if the user's password is correctly supplied.
# TYPE  DATABASE        USER            ADDRESS          METHOD
host    postgres        all             192.168.12.10/32
scram-sha-256

# Allow any user from hosts in the example.com domain to connect to
any database if the user's password is correctly supplied.
# Require SCRAM authentication for most users, but make an
exception for user 'mike', who uses an older client that doesn't support
SCRAM
# TYPE  DATABASE        USER            ADDRESS          METHOD
host    all             mike            .example.com
       all             all             .example.com
       .example.com
       md5
       .example.com
       scramble-sha-256

# In the absence of preceding "host" lines, these two lines will reject
all connections from 192.168.54.1 (since that entry will be matched first),
but allow GSSAPI connections from anywhere else on the Internet. The zero mask
causes no bits of the host IP address to be considered, so it matches any host.
# TYPE  DATABASE        USER            ADDRESS          METHOD
host    all             all             192.168.54.1/32
reject
host    all             all             0.0.0.0/0
       gss
```

600
# Allow users from 192.168.x.x hosts to connect to any database, if
# they pass the ident check. If, for example, ident says the user
# is
# "bryanh" and he requests to connect as PostgreSQL user "guest1",
# the
# connection is allowed if there is an entry in pg_ident.conf for
# map
# "omicron" that says "bryanh" is allowed to connect as "guest1".
# TYPE DATABASE USER ADDRESS
METHOD
host all all 192.168.0.0/16
ident map=omicron

# If these are the only three lines for local connections, they
# will
# allow local users to connect only to their own databases
(databases
# with the same name as their database user name) except for
# administrators
# and members of role "support", who can connect to all databases.
The file
# $PGDATA/admins contains a list of names of administrators.
Passwords
# are required in all cases.
# TYPE DATABASE USER ADDRESS
METHOD
local sameuser all md5
local all @admins md5
local all +support md5

# The last two lines above can be combined into a single line:
local all @admins,+support md5

# The database column can also use lists and file names:
local db1,db2,@demodbs all md5

## 20.2. User Name Maps

When using an external authentication system such as Ident or GSSAPI, the name of the operating
system user that initiated the connection might not be the same as the database user (role) that is to be
used. In this case, a user name map can be applied to map the operating system user name to a database
user. To use user name mapping, specify map=map-name in the options field in pg_hba.conf.
This option is supported for all authentication methods that receive external user names. Since different
mappings might be needed for different connections, the name of the map to be used is specified in the
map-name parameter in pg_hba.conf to indicate which map to use for each individual connection.

User name maps are defined in the ident map file, which by default is named pg_ident.conf and
is stored in the cluster's data directory. (It is possible to place the map file elsewhere, however; see the
ident_file configuration parameter.) The ident map file contains lines of the general form:

```
map-name system-username database-username
```

Comments and whitespace are handled in the same way as in pg_hba.conf. The map-name is
an arbitrary name that will be used to refer to this mapping in pg_hba.conf. The other two fields
Client Authentication

specify an operating system user name and a matching database user name. The same map-name can be used repeatedly to specify multiple user-mappings within a single map.

There is no restriction regarding how many database users a given operating system user can correspond to, nor vice versa. Thus, entries in a map should be thought of as meaning “this operating system user is allowed to connect as this database user”, rather than implying that they are equivalent. The connection will be allowed if there is any map entry that pairs the user name obtained from the external authentication system with the database user name that the user has requested to connect as.

If the system-username field starts with a slash (/), the remainder of the field is treated as a regular expression. (See Section 9.7.3.1 for details of PostgreSQL’s regular expression syntax.) The regular expression can include a single capture, or parenthesized subexpression, which can then be referenced in the database-username field as \1 (backslash-one). This allows the mapping of multiple user names in a single line, which is particularly useful for simple syntax substitutions. For example, these entries

mymap /(.*)@mydomain\.com$/ \1
mymap /(.*)@otherdomain\.com$/ guest

will remove the domain part for users with system user names that end with @mydomain.com, and allow any user whose system name ends with @otherdomain.com to log in as guest.

Tip
Keep in mind that by default, a regular expression can match just part of a string. It's usually wise to use ^ and $, as shown in the above example, to force the match to be to the entire system user name.

The pg_ident.conf file is read on start-up and when the main server process receives a SIGHUP signal. If you edit the file on an active system, you will need to signal the postmaster (using pg_ctl reload or kill -HUP) to make it re-read the file.

A pg_ident.conf file that could be used in conjunction with the pg_hba.conf file in Example 20.1 is shown in Example 20.2. In this example, anyone logged in to a machine on the 192.168 network that does not have the operating system user name bryanh, ann, or robert would not be granted access. Unix user robert would only be allowed access when he tries to connect as PostgreSQL user bob, not as robert or anyone else. ann would only be allowed to connect as ann. User bryanh would be allowed to connect as either bryanh or as guest1.

Example 20.2. An Example pg_ident.conf File

```
# MAPNAME       SYSTEM-USERNAME         PG-USERNAME
omicron         bryanh                  bryanh
omicron         ann                     ann
# bob has user name robert on these machines
omicron         robert                  bob
# bryanh can also connect as guest1
omicron         bryanh                  guest1
```

20.3. Authentication Methods

The following sections describe the authentication methods in more detail.
20.4. Trust Authentication

When trust authentication is specified, PostgreSQL assumes that anyone who can connect to the server is authorized to access the database with whatever database user name they specify (even superuser names). Of course, restrictions made in the database and user columns still apply. This method should only be used when there is adequate operating-system-level protection on connections to the server.

trust authentication is appropriate and very convenient for local connections on a single-user workstation. It is usually not appropriate by itself on a multiuser machine. However, you might be able to use trust even on a multiuser machine, if you restrict access to the server's Unix-domain socket file using file-system permissions. To do this, set the unix_socket_permissions (and possibly unix_socket_group) configuration parameters as described in Section 19.3. Or you could set the unix_socket_directories configuration parameter to place the socket file in a suitably restricted directory.

Setting file-system permissions only helps for Unix-socket connections. Local TCP/IP connections are not restricted by file-system permissions. Therefore, if you want to use file-system permissions for local security, remove the host ... 127.0.0.1 ... line from pg_hba.conf, or change it to a non-trust authentication method.

trust authentication is only suitable for TCP/IP connections if you trust every user on every machine that is allowed to connect to the server by the pg_hba.conf lines that specify trust. It is seldom reasonable to use trust for any TCP/IP connections other than those from localhost (127.0.0.1).

20.5. Password Authentication

There are several password-based authentication methods. These methods operate similarly but differ in how the users' passwords are stored on the server and how the password provided by a client is sent across the connection.

scram-sha-256

The method scram-sha-256 performs SCRAM-SHA-256 authentication, as described in RFC 7677. It is a challenge-response scheme that prevents password sniffing on untrusted connections and supports storing passwords on the server in a cryptographically hashed form that is thought to be secure.

This is the most secure of the currently provided methods, but it is not supported by older client libraries.

md5

The method md5 uses a custom less secure challenge-response mechanism. It prevents password sniffing and avoids storing passwords on the server in plain text but provides no protection if an attacker manages to steal the password hash from the server. Also, the MD5 hash algorithm is nowadays no longer considered secure against determined attacks.

The md5 method cannot be used with the db_user_namespace feature.

To ease transition from the md5 method to the newer SCRAM method, if md5 is specified as a method in pg_hba.conf but the user's password on the server is encrypted for SCRAM (see below), then SCRAM-based authentication will automatically be chosen instead.

password

The method password sends the password in clear-text and is therefore vulnerable to password "sniffing" attacks. It should always be avoided if possible. If the connection is protected by SSL

1 https://tools.ietf.org/html/rfc7677
encryption then password can be used safely, though. (Though SSL certificate authentication might be a better choice if one is depending on using SSL).

PostgreSQL database passwords are separate from operating system user passwords. The password for each database user is stored in the pg_authid system catalog. Passwords can be managed with the SQL commands CREATE ROLE and ALTER ROLE, e.g., `CREATE ROLE foo WITH LOGIN PASSWORD 'secret'`, or the psql command `\password`. If no password has been set up for a user, the stored password is null and password authentication will always fail for that user.

The availability of the different password-based authentication methods depends on how a user's password on the server is encrypted (or hashed, more accurately). This is controlled by the configuration parameter password_encryption at the time the password is set. If a password was encrypted using the `scram-sha-256` setting, then it can be used for the authentication methods `scram-sha-256` and `password` (but password transmission will be in plain text in the latter case). The authentication method specification `md5` will automatically switch to using the `scram-sha-256` method in this case, as explained above, so it will also work. If a password was encrypted using the `md5` setting, then it can be used only for the `md5` and `password` authentication method specifications (again, with the password transmitted in plain text in the latter case). (Previous PostgreSQL releases supported storing the password on the server in plain text. This is no longer possible.) To check the currently stored password hashes, see the system catalog `pg_authid`.

To upgrade an existing installation from `md5` to `scram-sha-256`, after having ensured that all client libraries in use are new enough to support SCRAM, set `password_encryption = 'scram-sha-256'` in `postgresql.conf`, make all users set new passwords, and change the authentication method specifications in `pg_hba.conf` to `scram-sha-256`.

## 20.6. GSSAPI Authentication

GSSAPI is an industry-standard protocol for secure authentication defined in RFC 2743. PostgreSQL supports GSSAPI with Kerberos authentication according to RFC 1964. GSSAPI provides automatic authentication (single sign-on) for systems that support it. The authentication itself is secure, but the data sent over the database connection will be sent unencrypted unless SSL is used.

GSSAPI support has to be enabled when PostgreSQL is built; see Chapter 16 for more information.

When GSSAPI uses Kerberos, it uses a standard principal in the format `servicename/host-name@realm`. The PostgreSQL server will accept any principal that is included in the keytab used by the server, but care needs to be taken to specify the correct principal details when making the connection from the client using the `krbsrvname` connection parameter. (See also Section 34.1.2.) The installation default can be changed from the default `SCRAM` at build time using `--with-krb-srvnam=whatever`. In most environments, this parameter never needs to be changed. Some Kerberos implementations might require a different service name, such as Microsoft Active Directory which requires the service name to be in upper case (`POSTGRES`).

`hostname` is the fully qualified host name of the server machine. The service principal's realm is the preferred realm of the server machine.

Client principals can be mapped to different PostgreSQL database user names with `pg_ident.conf`. For example, `pgusername@realm` could be mapped to just `pgusername`. Alternatively, you can use the full `username@realm` principal as the role name in PostgreSQL without any mapping.

PostgreSQL also supports a parameter to strip the realm from the principal. This method is supported for backwards compatibility and is strongly discouraged as it is then impossible to distinguish different users with the same user name but coming from different realms. To enable this, set `include_realm` to 0. For simple single-realm installations, doing that combined with setting the `krb_realm` parameter (which checks that the principal's realm matches exactly what is in the `krb_realm` parameter) is still secure; but this is a less capable approach compared to specifying an explicit mapping in `pg_ident.conf`.

604
Client Authentication

Make sure that your server keytab file is readable (and preferably only readable, not writable) by the PostgreSQL server account. (See also Section 18.1.) The location of the key file is specified by the \texttt{krb\_server\_keyfile} configuration parameter. The default is \texttt{/usr/local/pgsql/etc/krb5.keytab} (or whatever directory was specified as \texttt{sysconfdir} at build time). For security reasons, it is recommended to use a separate keytab just for the PostgreSQL server rather than opening up permissions on the system keytab file.

The keytab file is generated by the Kerberos software; see the Kerberos documentation for details. The following example is for MIT-compatible Kerberos 5 implementations:

\begin{verbatim}
  kadmin% ank -randkey postgres/server.my.domain.org
  kadmin% ktadd -k krb5.keytab postgres/server.my.domain.org
\end{verbatim}

When connecting to the database make sure you have a ticket for a principal matching the requested database user name. For example, for database user name \texttt{fred}, principal \texttt{fred@EXAMPLE.COM} would be able to connect. To also allow principal \texttt{fred/users.example.com@EXAMPLE.COM}, use a user name map, as described in Section 20.2.

The following configuration options are supported for GSSAPI:

\begin{description}
\item[include\_realm] If set to 0, the realm name from the authenticated user principal is stripped off before being passed through the user name mapping (Section 20.2). This is discouraged and is primarily available for backwards compatibility, as it is not secure in multi-realm environments unless \texttt{krb\_realm} is also used. It is recommended to leave \texttt{include\_realm} set to the default (1) and to provide an explicit mapping in \texttt{pg\_ident\_conf} to convert principal names to PostgreSQL user names.

\item[map] Allows for mapping between system and database user names. See Section 20.2 for details. For a GSSAPI/Kerberos principal, such as \texttt{username@EXAMPLE.COM} (or, less commonly, \texttt{username/hostbased@EXAMPLE.COM}), the user name used for mapping is \texttt{username@EXAMPLE.COM} (or \texttt{username/hostbased@EXAMPLE.COM}, respectively), unless \texttt{include\_realm} has been set to 0, in which case \texttt{username} (or \texttt{username/hostbased}) is what is seen as the system user name when mapping.

\item[krb\_realm] Sets the realm to match user principal names against. If this parameter is set, only users of that realm will be accepted. If it is not set, users of any realm can connect, subject to whatever user name mapping is done.
\end{description}

### 20.7. SSPI Authentication

SSPI is a Windows technology for secure authentication with single sign-on. PostgreSQL will use SSPI in \texttt{negotiate} mode, which will use Kerberos when possible and automatically fall back to NTLM in other cases. SSPI authentication only works when both server and client are running Windows, or, on non-Windows platforms, when GSSAPI is available.

When using Kerberos authentication, SSPI works the same way GSSAPI does; see Section 20.6 for details.

The following configuration options are supported for SSPI:

\begin{description}
\item[include\_realm] If set to 0, the realm name from the authenticated user principal is stripped off before being passed through the user name mapping (Section 20.2). This is discouraged and is primarily available for
backwards compatibility, as it is not secure in multi-realm environments unless krb_realm is also used. It is recommended to leave include_realm set to the default (1) and to provide an explicit mapping in `pg_ident.conf` to convert principal names to PostgreSQL user names.

**compat_realm**

If set to 1, the domain's SAM-compatible name (also known as the NetBIOS name) is used for the include_realm option. This is the default. If set to 0, the true realm name from the Kerberos user principal name is used.

Do not disable this option unless your server runs under a domain account (this includes virtual service accounts on a domain member system) and all clients authenticating through SSPI are also using domain accounts, or authentication will fail.

**upn_username**

If this option is enabled along with compat_realm, the user name from the Kerberos UPN is used for authentication. If it is disabled (the default), the SAM-compatible user name is used. By default, these two names are identical for new user accounts.

Note that libpq uses the SAM-compatible name if no explicit user name is specified. If you use libpq or a driver based on it, you should leave this option disabled or explicitly specify user name in the connection string.

**map**

Allows for mapping between system and database user names. See Section 20.2 for details. For a SSPI/Kerberos principal, such as `username@EXAMPLE.COM` (or, less commonly, `username/hostbased@EXAMPLE.COM`), the user name used for mapping is `username@EXAMPLE.COM` (or `username/hostbased@EXAMPLE.COM`, respectively), unless include_realm has been set to 0, in which case `username` (or `username/hostbased`) is what is seen as the system user name when mapping.

**krb_realm**

Sets the realm to match user principal names against. If this parameter is set, only users of that realm will be accepted. If it is not set, users of any realm can connect, subject to whatever user name mapping is done.

## 20.8. Ident Authentication

The ident authentication method works by obtaining the client's operating system user name from an ident server and using it as the allowed database user name (with an optional user name mapping). This is only supported on TCP/IP connections.

**Note**

When ident is specified for a local (non-TCP/IP) connection, peer authentication (see Section 20.9) will be used instead.

The following configuration options are supported for ident:

**map**

Allows for mapping between system and database user names. See Section 20.2 for details.

The “Identification Protocol” is described in RFC 1413. Virtually every Unix-like operating system ships with an ident server that listens on TCP port 113 by default. The basic functionality of an ident...
server is to answer questions like “What user initiated the connection that goes out of your port X and connects to my port Y?” Since PostgreSQL knows both X and Y when a physical connection is established, it can interrogate the ident server on the host of the connecting client and can theoretically determine the operating system user for any given connection.

The drawback of this procedure is that it depends on the integrity of the client: if the client machine is untrusted or compromised, an attacker could run just about any program on port 113 and return any user name they choose. This authentication method is therefore only appropriate for closed networks where each client machine is under tight control and where the database and system administrators operate in close contact. In other words, you must trust the machine running the ident server. Heed the warning:

The Identification Protocol is not intended as an authorization or access control protocol.

—RFC 1413

Some ident servers have a nonstandard option that causes the returned user name to be encrypted, using a key that only the originating machine's administrator knows. This option must not be used when using the ident server with PostgreSQL, since PostgreSQL does not have any way to decrypt the returned string to determine the actual user name.

20.9. Peer Authentication

The peer authentication method works by obtaining the client's operating system user name from the kernel and using it as the allowed database user name (with optional user name mapping). This method is only supported on local connections.

The following configuration options are supported for peer:

map

Allows for mapping between system and database user names. See Section 20.2 for details.

Peer authentication is only available on operating systems providing the getpeereid() function, the SO_PEERCRED socket parameter, or similar mechanisms. Currently that includes Linux, most flavors of BSD including macOS, and Solaris.

20.10. LDAP Authentication

This authentication method operates similarly to password except that it uses LDAP as the password verification method. LDAP is used only to validate the user name/password pairs. Therefore the user must already exist in the database before LDAP can be used for authentication.

LDAP authentication can operate in two modes. In the first mode, which we will call the simple bind mode, the server will bind to the distinguished name constructed as prefix username suffix. Typically, the prefix parameter is used to specify cn=, or DOMAIN\ in an Active Directory environment. suffix is used to specify the remaining part of the DN in a non-Active Directory environment.

In the second mode, which we will call the search+bind mode, the server first binds to the LDAP directory with a fixed user name and password, specified with ldapbinddn and ldapbindpasswd, and performs a search for the user trying to log in to the database. If no user and password is configured, an anonymous bind will be attempted to the directory. The search will be performed over the subtree at ldapbasedn, and will try to do an exact match of the attribute specified in ldapsearchattribute. Once the user has been found in this search, the server disconnects and re-binds to the directory as this user, using the password specified by the client, to verify that the login is correct. This mode is the same as that used by LDAP authentication schemes in other software, such as Apache.
mod_authnz_ldap and pam_ldap. This method allows for significantly more flexibility in where
the user objects are located in the directory, but will cause two separate connections to the LDAP
server to be made.

The following configuration options are used in both modes:

**ldapserver**

Names or IP addresses of LDAP servers to connect to. Multiple servers may be specified, sepa-
rated by spaces.

**ldapport**

Port number on LDAP server to connect to. If no port is specified, the LDAP library’s default
port setting will be used.

**ldapscheme**

Set to ldaps to use LDAPS. This is a non-standard way of using LDAP over SSL, supported by
some LDAP server implementations. See also the ldaptls option for an alternative.

**ldaptls**

Set to 1 to make the connection between PostgreSQL and the LDAP server use TLS encryption.
This uses the StartTLS operation per RFC 4513. See also the ldapscheme option for an
alternative.

Note that using ldapscheme or ldaptls only encrypts the traffic between the PostgreSQL server
and the LDAP server. The connection between the PostgreSQL server and the PostgreSQL client will
still be unencrypted unless SSL is used there as well.

The following options are used in simple bind mode only:

**ldapprefix**

String to prepend to the user name when forming the DN to bind as, when doing simple bind
authentication.

**ldapsuffix**

String to append to the user name when forming the DN to bind as, when doing simple bind
authentication.

The following options are used in search+bind mode only:

**ldapbasedn**

Root DN to begin the search for the user in, when doing search+bind authentication.

**ldapbinddn**

DN of user to bind to the directory with to perform the search when doing search+bind authen-
tication.

**ldapbindpasswd**

Password for user to bind to the directory with to perform the search when doing search+bind
authentication.

**ldapsearchattribute**

Attribute to match against the user name in the search when doing search+bind authentication. If
no attribute is specified, the uid attribute will be used.
**Client Authentication**

**ldapsearchfilter**

The search filter to use when doing search+bind authentication. Occurrences of $username will be replaced with the user name. This allows for more flexible search filters than ldapsearchattribute.

**ldapurl**

An RFC 4516 LDAP URL. This is an alternative way to write some of the other LDAP options in a more compact and standard form. The format is

```
ldap[s]://host[:port]/basedn?[?attribute]?[?scope]?[?filter]]
```

`scope` must be one of `base`, `one`, `sub`, typically the last. (The default is `base`, which is normally not useful in this application.) `attribute` can nominate a single attribute, in which case it is used as a value for `ldapsearchattribute`. If `attribute` is empty then `filter` can be used as a value for `ldapsearchfilter`.

The URL scheme `ldaps` chooses the LDAPS method for making LDAP connections over SSL, equivalent to using `ldapscheme=ldaps`. To use encrypted LDAP connections using the `StartTLS` operation, use the normal URL scheme `ldap` and specify the `ldaptls` option in addition to `ldapurl`.

For non-anonymous binds, `ldapbinddn` and `ldapbindpasswd` must be specified as separate options.

LDAP URLs are currently only supported with OpenLDAP, not on Windows.

It is an error to mix configuration options for simple bind with options for search+bind.

When using search+bind mode, the search can be performed using a single attribute specified with `ldapsearchattribute`, or using a custom search filter specified with `ldapsearchfilter`. Specifying `ldapsearchattribute=foo` is equivalent to specifying `ldapsearchfilter="(foo=$username)"`. If neither option is specified the default is `ldapsearchattribute=uid`.

Here is an example for a simple-bind LDAP configuration:

```
host ... ldap ldapscheme=ldaps
ldapprefix="cn="
ldapsuffix="", dc=example, dc=net"
```

When a connection to the database server as database user `someuser` is requested, PostgreSQL will attempt to bind to the LDAP server using the DN `cn=someuser, dc=example, dc=net` and the password provided by the client. If that connection succeeds, the database access is granted.

Here is an example for a search+bind configuration:

```
host ... ldap ldapscheme=ldaps
ldapprefix="dc=example, dc=net"
ldapsearchattribute=uid
```

When a connection to the database server as database user `someuser` is requested, PostgreSQL will attempt to bind anonymously (since `ldapbinddn` was not specified) to the LDAP server, perform a search for `(uid=someuser)` under the specified base DN. If an entry is found, it will then attempt to bind using that found information and the password supplied by the client. If that second connection succeeds, the database access is granted.

Here is the same search+bind configuration written as a URL:
Client Authentication

host ... ldap ldapurl="ldap://ldap.example.net/dc=example,dc=net?uid?sub"

Some other software that supports authentication against LDAP uses the same URL format, so it will be easier to share the configuration.

Here is an example for a search+bind configuration that uses ldapsearchfilter instead of ldapsearchattribute to allow authentication by user ID or email address:

host ... ldaps server=ldap.example.net ldaps baseddn="dc=example,dc=net" ldapsearchfilter="(|(uid=$username)(mail=$username))"

Tip
Since LDAP often uses commas and spaces to separate the different parts of a DN, it is often necessary to use double-quoted parameter values when configuring LDAP options, as shown in the examples.

20.11. RADIUS Authentication

This authentication method operates similarly to password except that it uses RADIUS as the password verification method. RADIUS is used only to validate the user name/password pairs. Therefore the user must already exist in the database before RADIUS can be used for authentication.

When using RADIUS authentication, an Access Request message will be sent to the configured RADIUS server. This request will be of type Authenticate Only, and include parameters for user name, password (encrypted) and NAS Identifier. The request will be encrypted using a secret shared with the server. The RADIUS server will respond to this server with either Access Accept or Access Reject. There is no support for RADIUS accounting.

Multiple RADIUS servers can be specified, in which case they will be tried sequentially. If a negative response is received from a server, the authentication will fail. If no response is received, the next server in the list will be tried. To specify multiple servers, put the names within quotes and separate the server names with a comma. If multiple servers are specified, all other RADIUS options can also be given as a comma separate list, to apply individual values to each server. They can also be specified as a single value, in which case this value will apply to all servers.

The following configuration options are supported for RADIUS:

**radiusservers**

The name or IP addresses of the RADIUS servers to connect to. This parameter is required.

**radiussecrets**

The shared secrets used when talking securely to the RADIUS server. This must have exactly the same value on the PostgreSQL and RADIUS servers. It is recommended that this be a string of at least 16 characters. This parameter is required.

Note

The encryption vector used will only be cryptographically strong if PostgreSQL is built with support for OpenSSL. In other cases, the transmission to the RADIUS server will not be secure.
Client Authentication

server should only be considered obfuscated, not secured, and external security measures should be applied if necessary.

radiusports

The port number on the RADIUS servers to connect to. If no port is specified, the default port 1812 will be used.

radiusidentifiers

The string used as NAS Identifier in the RADIUS requests. This parameter can be used as a second parameter identifying for example which database user the user is attempting to authenticate as, which can be used for policy matching on the RADIUS server. If no identifier is specified, the default postgresql will be used.


This authentication method uses SSL client certificates to perform authentication. It is therefore only available for SSL connections. When using this authentication method, the server will require that the client provide a valid, trusted certificate. No password prompt will be sent to the client. The cn (Common Name) attribute of the certificate will be compared to the requested database user name, and if they match the login will be allowed. User name mapping can be used to allow cn to be different from the database user name.

The following configuration options are supported for SSL certificate authentication:

map

Allows for mapping between system and database user names. See Section 20.2 for details.

In a pg_hba.conf record specifying certificate authentication, the authentication option clientcert is assumed to be 1, and it cannot be turned off since a client certificate is necessary for this method. What the cert method adds to the basic clientcert certificate validity test is a check that the cn attribute matches the database user name.

20.13. PAM Authentication

This authentication method operates similarly to password except that it uses PAM (Pluggable Authentication Modules) as the authentication mechanism. The default PAM service name is postgresql. PAM is used only to validate user name/password pairs and optionally the connected remote host name or IP address. Therefore the user must already exist in the database before PAM can be used for authentication. For more information about PAM, please read the Linux-PAM Page².

The following configuration options are supported for PAM:

pamservice

PAM service name.

pam_use_hostname

Determines whether the remote IP address or the host name is provided to PAM modules through the PAM_RHOST item. By default, the IP address is used. Set this option to 1 to use the resolved host name instead. Host name resolution can lead to login delays. (Most PAM configurations don’t use this information, so it is only necessary to consider this setting if a PAM configuration was specifically created to make use of it.)

² https://www.kernel.org/pub/linux/libs/pam/
20.14. BSD Authentication

This authentication method operates similarly to password except that it uses BSD Authentication to verify the password. BSD Authentication is used only to validate user name/password pairs. Therefore the user's role must already exist in the database before BSD Authentication can be used for authentication. The BSD Authentication framework is currently only available on OpenBSD.

BSD Authentication in PostgreSQL uses the auth-postgresql login type and authenticates with the postgresql login class if that's defined in login.conf. By default that login class does not exist, and PostgreSQL will use the default login class.

Note

To use BSD Authentication, the PostgreSQL user account (that is, the operating system user running the server) must first be added to the auth group. The auth group exists by default on OpenBSD systems.

20.15. Authentication Problems

Authentication failures and related problems generally manifest themselves through error messages like the following:

FATAL: no pg_hba.conf entry for host "123.123.123.123", user "andym", database "testdb"

This is what you are most likely to get if you succeed in contacting the server, but it does not want to talk to you. As the message suggests, the server refused the connection request because it found no matching entry in its pg_hba.conf configuration file.

FATAL: password authentication failed for user "andym"

Messages like this indicate that you contacted the server, and it is willing to talk to you, but not until you pass the authorization method specified in the pg_hba.conf file. Check the password you are providing, or check your Kerberos or ident software if the complaint mentions one of those authentication types.

FATAL: user "andym" does not exist

The indicated database user name was not found.

FATAL: database "testdb" does not exist

The database you are trying to connect to does not exist. Note that if you do not specify a database name, it defaults to the database user name, which might or might not be the right thing.
Tip

The server log might contain more information about an authentication failure than is reported to the client. If you are confused about the reason for a failure, check the server log.
Chapter 21. Database Roles

PostgreSQL manages database access permissions using the concept of *roles*. A role can be thought of as either a database user, or a group of database users, depending on how the role is set up. Roles can own database objects (for example, tables and functions) and can assign privileges on those objects to other roles to control who has access to which objects. Furthermore, it is possible to grant membership in a role to another role, thus allowing the member role to use privileges assigned to another role.

The concept of roles subsumes the concepts of “users” and “groups”. In PostgreSQL versions before 8.1, users and groups were distinct kinds of entities, but now there are only roles. Any role can act as a user, a group, or both.

This chapter describes how to create and manage roles. More information about the effects of role privileges on various database objects can be found in Section 5.6.

### 21.1. Database Roles

Database roles are conceptually completely separate from operating system users. In practice it might be convenient to maintain a correspondence, but this is not required. Database roles are global across a database cluster installation (and not per individual database). To create a role use the CREATE ROLE SQL command:

```sql
CREATE ROLE name;
```

`name` follows the rules for SQL identifiers: either unadorned without special characters, or double-quoted. (In practice, you will usually want to add additional options, such as LOGIN, to the command. More details appear below.) To remove an existing role, use the analogous DROP ROLE command:

```sql
DROP ROLE name;
```

For convenience, the programs createuser and dropuser are provided as wrappers around these SQL commands that can be called from the shell command line:

```bash
createuser name
dropuser name
```

To determine the set of existing roles, examine the `pg_roles` system catalog, for example

```sql
SELECT rolname FROM pg_roles;
```

The `psql` program's `\du` meta-command is also useful for listing the existing roles.

In order to bootstrap the database system, a freshly initialized system always contains one predefined role. This role is always a “superuser”, and by default (unless altered when running `initdb`) it will have the same name as the operating system user that initialized the database cluster. Customarily, this role will be named `postgres`. In order to create more roles you first have to connect as this initial role.

Every connection to the database server is made using the name of some particular role, and this role determines the initial access privileges for commands issued in that connection. The role name to use for a particular database connection is indicated by the client that is initiating the connection request in an application-specific fashion. For example, the `psql` program uses the `-u` command line option...
to indicate the role to connect as. Many applications assume the name of the current operating system user by default (including createuser and psql). Therefore it is often convenient to maintain a naming correspondence between roles and operating system users.

The set of database roles a given client connection can connect as is determined by the client authentication setup, as explained in Chapter 20. (Thus, a client is not limited to connect as the role matching its operating system user, just as a person's login name need not match his or her real name.) Since the role identity determines the set of privileges available to a connected client, it is important to carefully configure privileges when setting up a multiuser environment.

### 21.2. Role Attributes

A database role can have a number of attributes that define its privileges and interact with the client authentication system.

**login privilege**

Only roles that have the LOGIN attribute can be used as the initial role name for a database connection. A role with the LOGIN attribute can be considered the same as a "database user". To create a role with login privilege, use either:

- `CREATE ROLE name LOGIN;
- `CREATE USER name;`

(CREATE USER is equivalent to CREATE ROLE except that CREATE USER includes LOGIN by default, while CREATE ROLE does not.)

**superuser status**

A database superuser bypasses all permission checks, except the right to log in. This is a dangerous privilege and should not be used carelessly; it is best to do most of your work as a role that is not a superuser. To create a new database superuser, use `CREATE ROLE name SUPERUSER`. You must do this as a role that is already a superuser.

**database creation**

A role must be explicitly given permission to create databases (except for superusers, since those bypass all permission checks). To create such a role, use `CREATE ROLE name CREATEDB`.

**role creation**

A role must be explicitly given permission to create more roles (except for superusers, since those bypass all permission checks). To create such a role, use `CREATE ROLE name CREATEROLE`. A role with CREATEROLE privilege can alter and drop other roles, too, as well as grant or revoke membership in them. However, to create, alter, drop, or change membership of a superuser role, superuser status is required; CREATEROLE is insufficient for that.

**initiating replication**

A role must explicitly be given permission to initiate streaming replication (except for superusers, since those bypass all permission checks). A role used for streaming replication must have LOGIN permission as well. To create such a role, use `CREATE ROLE name REPLICATION LOGIN`.

**password**

A password is only significant if the client authentication method requires the user to supply a password when connecting to the database. The password and md5 authentication methods make use of passwords. Database passwords are separate from operating system passwords. Specify a password upon role creation with `CREATE ROLE name PASSWORD 'string'`. 

615
A role's attributes can be modified after creation with ALTER ROLE. See the reference pages for the CREATE ROLE and ALTER ROLE commands for details.

### Tip

It is good practice to create a role that has the CREATEDB and CREATEROLE privileges, but is not a superuser, and then use this role for all routine management of databases and roles. This approach avoids the dangers of operating as a superuser for tasks that do not really require it.

A role can also have role-specific defaults for many of the run-time configuration settings described in Chapter 19. For example, if for some reason you want to disable index scans (hint: not a good idea) anytime you connect, you can use:

```sql
ALTER ROLE myname SET enable_indexscan TO off;
```

This will save the setting (but not set it immediately). In subsequent connections by this role it will appear as though `SET enable_indexscan TO off` had been executed just before the session started. You can still alter this setting during the session; it will only be the default. To remove a role-specific default setting, use `ALTER ROLE rolename RESET varname`. Note that role-specific defaults attached to roles without LOGIN privilege are fairly useless, since they will never be invoked.

### 21.3. Role Membership

It is frequently convenient to group users together to ease management of privileges: that way, privileges can be granted to, or revoked from, a group as a whole. In PostgreSQL this is done by creating a role that represents the group, and then granting *membership* in the group role to individual user roles.

To set up a group role, first create the role:

```sql
CREATE ROLE name;
```

Typically a role being used as a group would not have the LOGIN attribute, though you can set it if you wish.

Once the group role exists, you can add and remove members using the GRANT and REVOKE commands:

```sql
GRANT group_role TO role1, ... ;
REVOKE group_role FROM role1, ... ;
```

You can grant membership to other group roles, too (since there isn't really any distinction between group roles and non-group roles). The database will not let you set up circular membership loops. Also, it is not permitted to grant membership in a role to PUBLIC.

The members of a group role can use the privileges of the role in two ways. First, every member of a group can explicitly do SET ROLE to temporarily “become” the group role. In this state, the database session has access to the privileges of the group role rather than the original login role, and any database objects created are considered owned by the group role not the login role. Second, member roles that have the INHERIT attribute automatically have use of the privileges of roles of which they are members, including any privileges inherited by those roles. As an example, suppose we have done:

```sql
CREATE ROLE joe LOGIN INHERIT;
CREATE ROLE admin NOINHERIT;
```
CREATE ROLE wheel NOINHERIT;
GRANT admin TO joe;
GRANT wheel TO admin;

Immediately after connecting as role joe, a database session will have use of privileges granted directly to joe plus any privileges granted to admin, because joe “inherits” admin’s privileges. However, privileges granted to wheel are not available, because even though joe is indirectly a member of wheel, the membership is via admin which has the NOINHERIT attribute. After:

SET ROLE admin;

the session would have use of only those privileges granted to admin, and not those granted to joe. After:

SET ROLE wheel;

the session would have use of only those privileges granted to wheel, and not those granted to either joe or admin. The original privilege state can be restored with any of:

SET ROLE joe;
SET ROLE NONE;
RESET ROLE;

Note

The SET ROLE command always allows selecting any role that the original login role is directly or indirectly a member of. Thus, in the above example, it is not necessary to become admin before becoming wheel.

Note

In the SQL standard, there is a clear distinction between users and roles, and users do not automatically inherit privileges while roles do. This behavior can be obtained in PostgreSQL by giving roles being used as SQL roles the INHERIT attribute, while giving roles being used as SQL users the NOINHERIT attribute. However, PostgreSQL defaults to giving all roles the INHERIT attribute, for backward compatibility with pre-8.1 releases in which users always had use of permissions granted to groups they were members of.

The role attributes LOGIN, SUPERUSER, CREATEDB, and CREATEROLE can be thought of as special privileges, but they are never inherited as ordinary privileges on database objects are. You must actually SET ROLE to a specific role having one of these attributes in order to make use of the attribute. Continuing the above example, we might choose to grant CREATEDB and CREATEROLE to the admin role. Then a session connecting as role joe would not have these privileges immediately, only after doing SET ROLE admin.

To destroy a group role, use DROP ROLE:

DROP ROLE name;

Any memberships in the group role are automatically revoked (but the member roles are not otherwise affected).
21.4. Dropping Roles

Because roles can own database objects and can hold privileges to access other objects, dropping a role is often not just a matter of a quick DROP ROLE. Any objects owned by the role must first be dropped or reassigned to other owners; and any permissions granted to the role must be revoked.

Ownership of objects can be transferred one at a time using ALTER commands, for example:

```
ALTER TABLE bobs_table OWNER TO alice;
```

Alternatively, the REASSIGN OWNED command can be used to reassign ownership of all objects owned by the role-to-be-dropped to a single other role. Because REASSIGN OWNED cannot access objects in other databases, it is necessary to run it in each database that contains objects owned by the role. (Note that the first such REASSIGN OWNED will change the ownership of any shared-across-databases objects, that is databases or tablespaces, that are owned by the role-to-be-dropped.)

Once any valuable objects have been transferred to new owners, any remaining objects owned by the role-to-be-dropped can be dropped with the DROP OWNED command. Again, this command cannot access objects in other databases, so it is necessary to run it in each database that contains objects owned by the role. Also, DROP OWNED will not drop entire databases or tablespaces, so it is necessary to do that manually if the role owns any databases or tablespaces that have not been transferred to new owners.

DROP OWNED also takes care of removing any privileges granted to the target role for objects that do not belong to it. Because REASSIGN OWNED does not touch such objects, it's typically necessary to run both REASSIGN OWNED and DROP OWNED (in that order!) to fully remove the dependencies of a role to be dropped.

In short then, the most general recipe for removing a role that has been used to own objects is:

```
REASSIGN OWNED BY doomed_role TO successor_role;
DROP OWNED BY doomed_role;
-- repeat the above commands in each database of the cluster
DROP ROLE doomed_role;
```

When not all owned objects are to be transferred to the same successor owner, it's best to handle the exceptions manually and then perform the above steps to mop up.

If DROP ROLE is attempted while dependent objects still remain, it will issue messages identifying which objects need to be reassigned or dropped.

21.5. Default Roles

PostgreSQL provides a set of default roles which provide access to certain, commonly needed, privileged capabilities and information. Administrators can GRANT these roles to users and/or other roles in their environment, providing those users with access to the specified capabilities and information.

The default roles are described in Table 21.1. Note that the specific permissions for each of the default roles may change in the future as additional capabilities are added. Administrators should monitor the release notes for changes.

<table>
<thead>
<tr>
<th>Role</th>
<th>Allowed Access</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_read_all_settings</td>
<td>Read all configuration variables, even those normally visible only to superusers.</td>
</tr>
</tbody>
</table>
### Database Roles

<table>
<thead>
<tr>
<th>Role</th>
<th>Allowed Access</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_read_all_stats</td>
<td>Read all pg_stat_* views and use various statistics related extensions, even those normally visible only to superusers.</td>
</tr>
<tr>
<td>pg_stat_scan_tables</td>
<td>Execute monitoring functions that may take ACCESS SHARE locks on tables, potentially for a long time.</td>
</tr>
<tr>
<td>pg_signal_backend</td>
<td>Send signals to other backends (eg: cancel query, terminate).</td>
</tr>
<tr>
<td>pg_read_server_files</td>
<td>Allow reading files from any location the database can access on the server with COPY and other file-access functions.</td>
</tr>
<tr>
<td>pg_write_server_files</td>
<td>Allow writing to files in any location the database can access on the server with COPY and other file-access functions.</td>
</tr>
<tr>
<td>pg_execute_server_program</td>
<td>Allow executing programs on the database server as the user the database runs as with COPY and other functions which allow executing a server-side program.</td>
</tr>
<tr>
<td>pg_monitor</td>
<td>Read/execute various monitoring views and functions. This role is a member of pg_read_all_settings, pg_read_all_stats and pg_stat_scan_tables.</td>
</tr>
</tbody>
</table>

The pg_read_server_files, pg_write_server_files and pg_execute_server_program roles are intended to allow administrators to have trusted, but non-superuser, roles which are able to access files and run programs on the database server as the user the database runs as. As these roles are able to access any file on the server file system, they bypass all database-level permission checks when accessing files directly and they could be used to gain superuser-level access, therefore care should be taken when granting these roles to users.

The pg_monitor, pg_read_all_settings, pg_read_all_stats and pg_stat_scan_tables roles are intended to allow administrators to easily configure a role for the purpose of monitoring the database server. They grant a set of common privileges allowing the role to read various useful configuration settings, statistics and other system information normally restricted to superusers.

Care should be taken when granting these roles to ensure they are only used where needed and with the understanding that these roles grant access to privileged information.

Administrators can grant access to these roles to users using the GRANT command:

```
GRANT pg_signal_backend TO admin_user;
```

## 21.6. Function Security

Functions, triggers and row-level security policies allow users to insert code into the backend server that other users might execute unintentionally. Hence, these mechanisms permit users to “Trojan horse” others with relative ease. The strongest protection is tight control over who can define objects. Where that is infeasible, write queries referring only to objects having trusted owners. Remove from search_path the public schema and any other schemas that permit untrusted users to create objects.

Functions run inside the backend server process with the operating system permissions of the database server daemon. If the programming language used for the function allows unchecked memory accesses, it is possible to change the server’s internal data structures. Hence, among many other things, such
functions can circumvent any system access controls. Function languages that allow such access are considered “untrusted”, and PostgreSQL allows only superusers to create functions written in those languages.
Chapter 22. Managing Databases

Every instance of a running PostgreSQL server manages one or more databases. Databases are therefore the topmost hierarchical level for organizing SQL objects (“database objects”). This chapter describes the properties of databases, and how to create, manage, and destroy them.

22.1. Overview

A database is a named collection of SQL objects (“database objects”). Generally, every database object (tables, functions, etc.) belongs to one and only one database. (However there are a few system catalogs, for example \texttt{pg\_database}, that belong to a whole cluster and are accessible from each database within the cluster.) More accurately, a database is a collection of schemas and the schemas contain the tables, functions, etc. So the full hierarchy is: server, database, schema, table (or some other kind of object, such as a function).

When connecting to the database server, a client must specify in its connection request the name of the database it wants to connect to. It is not possible to access more than one database per connection. However, an application is not restricted in the number of connections it opens to the same or other databases. Databases are physically separated and access control is managed at the connection level. If one PostgreSQL server instance is to house projects or users that should be separate and for the most part unaware of each other, it is therefore recommended to put them into separate databases. If the projects or users are interrelated and should be able to use each other's resources, they should be put in the same database but possibly into separate schemas. Schemas are a purely logical structure and who can access what is managed by the privilege system. More information about managing schemas is in Section 5.8.

Databases are created with the \texttt{CREATE DATABASE} command (see Section 22.2) and destroyed with the \texttt{DROP DATABASE} command (see Section 22.5). To determine the set of existing databases, examine the \texttt{pg\_database} system catalog, for example

\begin{verbatim}
SELECT datname FROM pg_database;
\end{verbatim}

The psql program's \texttt{\textasciitilde l} meta-command and \texttt{-l} command-line option are also useful for listing the existing databases.

\begin{center}
\textbf{Note}

The SQL standard calls databases “catalogs”, but there is no difference in practice.
\end{center}

22.2. Creating a Database

In order to create a database, the PostgreSQL server must be up and running (see Section 18.3).

Databases are created with the SQL command \texttt{CREATE DATABASE}:

\begin{verbatim}
CREATE DATABASE name;
\end{verbatim}

where \texttt{name} follows the usual rules for SQL identifiers. The current role automatically becomes the owner of the new database. It is the privilege of the owner of a database to remove it later (which also removes all the objects in it, even if they have a different owner).

The creation of databases is a restricted operation. See Section 21.2 for how to grant permission.
Managing Databases

Since you need to be connected to the database server in order to execute the `CREATE DATABASE` command, the question remains how the first database at any given site can be created. The first database is always created by the `initdb` command when the data storage area is initialized. (See Section 18.2.) This database is called `postgres`. So to create the first “ordinary” database you can connect to `postgres`.

A second database, `template1`, is also created during database cluster initialization. Whenever a new database is created within the cluster, `template1` is essentially cloned. This means that any changes you make in `template1` are propagated to all subsequently created databases. Because of this, avoid creating objects in `template1` unless you want them propagated to every newly created database. More details appear in Section 22.3.

As a convenience, there is a program you can execute from the shell to create new databases, `createdb`.

`createdb dbname`

`createdb` does no magic. It connects to the `postgres` database and issues the `CREATE DATABASE` command, exactly as described above. The `createdb` reference page contains the invocation details. Note that `createdb` without any arguments will create a database with the current user name.

Note

Chapter 20 contains information about how to restrict who can connect to a given database.

Sometimes you want to create a database for someone else, and have them become the owner of the new database, so they can configure and manage it themselves. To achieve that, use one of the following commands:

```
CREATE DATABASE dbname OWNER rolename;
```

from the SQL environment, or:

```
createdb -O rolename dbname
```

from the shell. Only the superuser is allowed to create a database for someone else (that is, for a role you are not a member of).

22.3. Template Databases

`CREATE DATABASE` actually works by copying an existing database. By default, it copies the standard system database named `template1`. Thus that database is the “template” from which new databases are made. If you add objects to `template1`, these objects will be copied into subsequently created user databases. This behavior allows site-local modifications to the standard set of objects in databases. For example, if you install the procedural language PL/Perl in `template1`, it will automatically be available in user databases without any extra action being taken when those databases are created.

There is a second standard system database named `template0`. This database contains the same data as the initial contents of `template1`, that is, only the standard objects predefined by your version of PostgreSQL. `template0` should never be changed after the database cluster has been initialized. By instructing `CREATE DATABASE` to copy `template0` instead of `template1`, you can create a “virgin” user database that contains none of the site-local additions in `template1`. This is particularly handy when restoring a `pg_dump` dump: the dump script should be restored in a virgin database.
to ensure that one recreates the correct contents of the dumped database, without conflicting with objects that might have been added to template1 later on.

Another common reason for copying template0 instead of template1 is that new encoding and locale settings can be specified when copying template0, whereas a copy of template1 must use the same settings it does. This is because template1 might contain encoding-specific or locale-specific data, while template0 is known not to.

To create a database by copying template0, use:

```
CREATE DATABASE dbname TEMPLATE template0;
```

from the SQL environment, or:

```
createdb -T template0 dbname
```

from the shell.

It is possible to create additional template databases, and indeed one can copy any database in a cluster by specifying its name as the template for CREATE DATABASE. It is important to understand, however, that this is not (yet) intended as a general-purpose "COPY DATABASE" facility. The principal limitation is that no other sessions can be connected to the source database while it is being copied. CREATE DATABASE will fail if any other connection exists when it starts; during the copy operation, new connections to the source database are prevented.

Two useful flags exist in pg_database for each database: the columns datistemplate and datallowconn. datistemplate can be set to indicate that a database is intended as a template for CREATE DATABASE. If this flag is set, the database can be cloned by any user with CREATEDB privileges; if it is not set, only superusers and the owner of the database can clone it. If datallowconn is false, then no new connections to that database will be allowed (but existing sessions are not terminated simply by setting the flag false). The template0 database is normally marked datallowconn = false to prevent its modification. Both template0 and template1 should always be marked with datistemplate = true.

Note

template1 and template0 do not have any special status beyond the fact that the name template1 is the default source database name for CREATE DATABASE. For example, one could drop template1 and recreate it from template0 without any ill effects. This course of action might be advisable if one has carelessly added a bunch of junk in template1. (To delete template1, it must have pg_database.datistemplate = false.)

The postgres database is also created when a database cluster is initialized. This database is meant as a default database for users and applications to connect to. It is simply a copy of template1 and can be dropped and recreated if necessary.

22.4. Database Configuration

Recall from Chapter 19 that the PostgreSQL server provides a large number of run-time configuration variables. You can set database-specific default values for many of these settings.

For example, if for some reason you want to disable the GEQO optimizer for a given database, you’d ordinarily have to either disable it for all databases or make sure that every connecting client is careful to issue `SET geqo TO off`. To make this setting the default within a particular database, you can execute the command:
ALTER DATABASE mydb SET geqo TO off;

This will save the setting (but not set it immediately). In subsequent connections to this database it will appear as though SET geqo TO off; had been executed just before the session started. Note that users can still alter this setting during their sessions; it will only be the default. To undo any such setting, use ALTER DATABASE dbname RESET varname.

22.5. Destroying a Database

Databases are destroyed with the command DROP DATABASE:

DROP DATABASE name;

Only the owner of the database, or a superuser, can drop a database. Dropping a database removes all objects that were contained within the database. The destruction of a database cannot be undone.

You cannot execute the DROP DATABASE command while connected to the victim database. You can, however, be connected to any other database, including the template1 database. template1 would be the only option for dropping the last user database of a given cluster.

For convenience, there is also a shell program to drop databases, dropdb:

dropdb dbname

(Unlike createdb, it is not the default action to drop the database with the current user name.)

22.6. Tablespaces

Tablespaces in PostgreSQL allow database administrators to define locations in the file system where the files representing database objects can be stored. Once created, a tablespace can be referred to by name when creating database objects.

By using tablespaces, an administrator can control the disk layout of a PostgreSQL installation. This is useful in at least two ways. First, if the partition or volume on which the cluster was initialized runs out of space and cannot be extended, a tablespace can be created on a different partition and used until the system can be reconfigured.

Second, tablespaces allow an administrator to use knowledge of the usage pattern of database objects to optimize performance. For example, an index which is very heavily used can be placed on a very fast, highly available disk, such as an expensive solid state device. At the same time a table storing archived data which is rarely used or not performance critical could be stored on a less expensive, slower disk system.

**Warning**

Even though located outside the main PostgreSQL data directory, tablespaces are an integral part of the database cluster and cannot be treated as an autonomous collection of data files. They are dependent on metadata contained in the main data directory, and therefore cannot be attached to a different database cluster or backed up individually. Similarly, if you lose a tablespace (file deletion, disk failure, etc), the database cluster might become unreadable or unable to start. Placing a tablespace on a temporary file system like a RAM disk risks the reliability of the entire cluster.

To define a tablespace, use the CREATE TABLESPACE command, for example:


CREATE TABLESPACE fastspace LOCATION '/ssd1/postgresql/data';

The location must be an existing, empty directory that is owned by the PostgreSQL operating system user. All objects subsequently created within the tablespace will be stored in files underneath this directory. The location must not be on removable or transient storage, as the cluster might fail to function if the tablespace is missing or lost.

**Note**

There is usually not much point in making more than one tablespace per logical file system, since you cannot control the location of individual files within a logical file system. However, PostgreSQL does not enforce any such limitation, and indeed it is not directly aware of the file system boundaries on your system. It just stores files in the directories you tell it to use.

Creation of the tablespace itself must be done as a database superuser, but after that you can allow ordinary database users to use it. To do that, grant them the CREATE privilege on it.

Tables, indexes, and entire databases can be assigned to particular tablespaces. To do so, a user with the CREATE privilege on a given tablespace must pass the tablespace name as a parameter to the relevant command. For example, the following creates a table in the tablespace `space1`:

```
CREATE TABLE foo(i int) TABLESPACE space1;
```

Alternatively, use the default_tablespace parameter:

```
SET default_tablespace = space1;
CREATE TABLE foo(i int);
```

When `default_tablespace` is set to anything but an empty string, it supplies an implicit TABLESPACE clause for CREATE TABLE and CREATE INDEX commands that do not have an explicit one.

There is also a temp_tablespaces parameter, which determines the placement of temporary tables and indexes, as well as temporary files that are used for purposes such as sorting large data sets. This can be a list of tablespace names, rather than only one, so that the load associated with temporary objects can be spread over multiple tablespaces. A random member of the list is picked each time a temporary object is to be created.

The tablespace associated with a database is used to store the system catalogs of that database. Furthermore, it is the default tablespace used for tables, indexes, and temporary files created within the database, if no TABLESPACE clause is given and no other selection is specified by default_tablespace or temp_tablespaces (as appropriate). If a database is created without specifying a tablespace for it, it uses the same tablespace as the template database it is copied from.

Two tablespaces are automatically created when the database cluster is initialized. The `pg_global` tablespace is used for shared system catalogs. The `pg_default` tablespace is the default tablespace of the `template1` and `template0` databases (and, therefore, will be the default tablespace for other databases as well, unless overridden by a TABLESPACE clause in CREATE DATABASE).

Once created, a tablespace can be used from any database, provided the requesting user has sufficient privilege. This means that a tablespace cannot be dropped until all objects in all databases using the tablespace have been removed.

To remove an empty tablespace, use the DROP TABLESPACE command.
To determine the set of existing tablespaces, examine the `pg_tablespace` system catalog, for example

```sql
SELECT spcname FROM pg_tablespace;
```

The `psql` program's `\db` meta-command is also useful for listing the existing tablespaces.

PostgreSQL makes use of symbolic links to simplify the implementation of tablespaces. This means that tablespaces can be used only on systems that support symbolic links.

The directory `$PGDATA/pg_tblspc` contains symbolic links that point to each of the non-built-in tablespaces defined in the cluster. Although not recommended, it is possible to adjust the tablespace layout by hand by redefining these links. Under no circumstances perform this operation while the server is running. Note that in PostgreSQL 9.1 and earlier you will also need to update the `pg_tablespace` catalog with the new locations. (If you do not, `pg_dump` will continue to output the old tablespace locations.)
Chapter 23. Localization

This chapter describes the available localization features from the point of view of the administrator. PostgreSQL supports two localization facilities:

- Using the locale features of the operating system to provide locale-specific collation order, number formatting, translated messages, and other aspects. This is covered in Section 23.1 and Section 23.2.

- Providing a number of different character sets to support storing text in all kinds of languages, and providing character set translation between client and server. This is covered in Section 23.3.

23.1. Locale Support

Locale support refers to an application respecting cultural preferences regarding alphabets, sorting, number formatting, etc. PostgreSQL uses the standard ISO C and POSIX locale facilities provided by the server operating system. For additional information refer to the documentation of your system.

23.1.1. Overview

Locale support is automatically initialized when a database cluster is created using `initdb`. `initdb` will initialize the database cluster with the locale setting of its execution environment by default, so if your system is already set to use the locale that you want in your database cluster then there is nothing else you need to do. If you want to use a different locale (or you are not sure which locale your system is set to), you can instruct `initdb` exactly which locale to use by specifying the `--locale` option. For example:

```
initdb --locale=sv_SE
```

This example for Unix systems sets the locale to Swedish (sv) as spoken in Sweden (SE). Other possibilities might include `en_US` (U.S. English) and `fr_CA` (French Canadian). If more than one character set can be used for a locale then the specifications can take the form `language_territory.codeset`. For example, `fr_BE.UTF-8` represents the French language (fr) as spoken in Belgium (BE), with a UTF-8 character set encoding.

What locales are available on your system under what names depends on what was provided by the operating system vendor and what was installed. On most Unix systems, the command `locale -a` will provide a list of available locales. Windows uses more verbose locale names, such as `German_Germany` or `Swedish_Sweden.1252`, but the principles are the same.

Occasionally it is useful to mix rules from several locales, e.g., use English collation rules but Spanish messages. To support that, a set of locale subcategories exist that control only certain aspects of the localization rules:

- **LC_COLLATE**: String sort order
- **LC_CTYPE**: Character classification (What is a letter? Its upper-case equivalent?)
- **LC_MESSAGES**: Language of messages
- **LC_MONETARY**: Formatting of currency amounts
- **LC_NUMERIC**: Formatting of numbers
- **LC_TIME**: Formatting of dates and times

The category names translate into names of `initdb` options to override the locale choice for a specific category. For instance, to set the locale to French Canadian, but use U.S. rules for formatting currency, use `initdb --locale=fr_CA --lc-monetary=en_US`.
If you want the system to behave as if it had no locale support, use the special locale name C, or equivalently POSIX.

Some locale categories must have their values fixed when the database is created. You can use different settings for different databases, but once a database is created, you cannot change them for that database anymore. LC_COLLATE and LC_CTYPE are these categories. They affect the sort order of indexes, so they must be kept fixed, or indexes on text columns would become corrupt. (But you can alleviate this restriction using collations, as discussed in Section 23.2.) The default values for these categories are determined when initdb is run, and those values are used when new databases are created, unless specified otherwise in the CREATE DATABASE command.

The other locale categories can be changed whenever desired by setting the server configuration parameters that have the same name as the locale categories (see Section 19.11.2 for details). The values that are chosen by initdb are actually only written into the configuration file postgresql.conf to serve as defaults when the server is started. If you remove these assignments from postgresql.conf then the server will inherit the settings from its execution environment.

Note that the locale behavior of the server is determined by the environment variables seen by the server, not by the environment of any client. Therefore, be careful to configure the correct locale settings before starting the server. A consequence of this is that if client and server are set up in different locales, messages might appear in different languages depending on where they originated.

---

When we speak of inheriting the locale from the execution environment, this means the following on most operating systems: For a given locale category, say the collation, the following environment variables are consulted in this order until one is found to be set: LC_ALL, LC_COLLATE (or the variable corresponding to the respective category), LANG. If none of these environment variables are set then the locale defaults to C.

Some message localization libraries also look at the environment variable LANGUAGE which overrides all other locale settings for the purpose of setting the language of messages. If in doubt, please refer to the documentation of your operating system, in particular the documentation about gettext.

To enable messages to be translated to the user's preferred language, NLS must have been selected at build time (configure --enable-nls). All other locale support is built in automatically.

23.1.2. Behavior

The locale settings influence the following SQL features:

- Sort order in queries using ORDER BY or the standard comparison operators on textual data
- The upper, lower, and initcap functions
- Pattern matching operators (LIKE, SIMILAR TO, and POSIX-style regular expressions); locales affect both case insensitive matching and the classification of characters by character-class regular expressions
- The to_char family of functions
- The ability to use indexes with LIKE clauses

The drawback of using locales other than C or POSIX in PostgreSQL is its performance impact. It slows character handling and prevents ordinary indexes from being used by LIKE. For this reason use locales only if you actually need them.
As a workaround to allow PostgreSQL to use indexes with LIKE clauses under a non-C locale, several custom operator classes exist. These allow the creation of an index that performs a strict character-by-character comparison, ignoring locale comparison rules. Refer to Section 11.9 for more information. Another approach is to create indexes using the C collation, as discussed in Section 23.2.

### 23.1.3. Problems

If locale support doesn't work according to the explanation above, check that the locale support in your operating system is correctly configured. To check what locales are installed on your system, you can use the command `locale -a` if your operating system provides it.

Check that PostgreSQL is actually using the locale that you think it is. The LC_COLLATE and LC_CTYPE settings are determined when a database is created, and cannot be changed except by creating a new database. Other locale settings including LC_MESSAGES and LC_MONETARY are initially determined by the environment the server is started in, but can be changed on-the-fly. You can check the active locale settings using the `SHOW` command.

The directory `src/test/locale` in the source distribution contains a test suite for PostgreSQL's locale support.

Client applications that handle server-side errors by parsing the text of the error message will obviously have problems when the server's messages are in a different language. Authors of such applications are advised to make use of the error code scheme instead.

Maintaining catalogs of message translations requires the on-going efforts of many volunteers that want to see PostgreSQL speak their preferred language well. If messages in your language are currently not available or not fully translated, your assistance would be appreciated. If you want to help, refer to Chapter 55 or write to the developers' mailing list.

### 23.2. Collation Support

The collation feature allows specifying the sort order and character classification behavior of data per-column, or even per-operation. This alleviates the restriction that the LC_COLLATE and LC_CTYPE settings of a database cannot be changed after its creation.

#### 23.2.1. Concepts

Conceptually, every expression of a collatable data type has a collation. (The built-in collatable data types are `text`, `varchar`, and `char`. User-defined base types can also be marked collatable, and of course a domain over a collatable data type is collatable.) If the expression is a column reference, the collation of the expression is the defined collation of the column. If the expression is a constant, the collation is the default collation of the data type of the constant. The collation of a more complex expression is derived from the collations of its inputs, as described below.

The collation of an expression can be the "default" collation, which means the locale settings defined for the database. It is also possible for an expression's collation to be indeterminate. In such cases, ordering operations and other operations that need to know the collation will fail.

When the database system has to perform an ordering or a character classification, it uses the collation of the input expression. This happens, for example, with ORDER BY clauses and function or operator calls such as `<`. The collation to apply for an ORDER BY clause is simply the collation of the sort key. The collation to apply for a function or operator call is derived from the arguments, as described below. In addition to comparison operators, collations are taken into account by functions that convert between lower and upper case letters, such as `lower`, `upper`, and `initcap`; by pattern matching operators; and by `to_char` and related functions.

For a function or operator call, the collation that is derived by examining the argument collations is used at run time for performing the specified operation. If the result of the function or operator call is of
a collatable data type, the collation is also used at parse time as the defined collation of the function or operator expression, in case there is a surrounding expression that requires knowledge of its collation.

The **collation derivation** of an expression can be implicit or explicit. This distinction affects how collations are combined when multiple different collations appear in an expression. An explicit collation derivation occurs when a `COLLATE` clause is used; all other collation derivations are implicit. When multiple collations need to be combined, for example in a function call, the following rules are used:

1. If any input expression has an explicit collation derivation, then all explicitly derived collations among the input expressions must be the same, otherwise an error is raised. If any explicitly derived collation is present, that is the result of the collation combination.

2. Otherwise, all input expressions must have the same implicit collation derivation or the default collation. If any non-default collation is present, that is the result of the collation combination. Otherwise, the result is the default collation.

3. If there are conflicting non-default implicit collations among the input expressions, then the combination is deemed to have indeterminate collation. This is not an error condition unless the particular function being invoked requires knowledge of the collation it should apply. If it does, an error will be raised at run-time.

For example, consider this table definition:

```sql
CREATE TABLE test1 (
    a text COLLATE "de_DE",
    b text COLLATE "es_ES",
    ...,
);
```

Then in

```sql
SELECT a < 'foo' FROM test1;
```

the `<` comparison is performed according to `de_DE` rules, because the expression combines an implicitly derived collation with the default collation. But in

```sql
SELECT a < ('foo' COLLATE "fr_FR") FROM test1;
```

the comparison is performed using `fr_FR` rules, because the explicit collation derivation overrides the implicit one. Furthermore, given

```sql
SELECT a < b FROM test1;
```

the parser cannot determine which collation to apply, since the `a` and `b` columns have conflicting implicit collations. Since the `<` operator does need to know which collation to use, this will result in an error. The error can be resolved by attaching an explicit collation specifier to either input expression, thus:

```sql
SELECT a < b COLLATE "de_DE" FROM test1;
```

or equivalently

```sql
SELECT a COLLATE "de_DE" < b FROM test1;
```

On the other hand, the structurally similar case
Localization

SELECT a || b FROM test1;

does not result in an error, because the || operator does not care about collations: its result is the same regardless of the collation.

The collation assigned to a function or operator's combined input expressions is also considered to apply to the function or operator's result, if the function or operator delivers a result of a collatable data type. So, in

SELECT * FROM test1 ORDER BY a || 'foo';

the ordering will be done according to de_DE rules. But this query:

SELECT * FROM test1 ORDER BY a || b;

results in an error, because even though the || operator doesn't need to know a collation, the ORDER BY clause does. As before, the conflict can be resolved with an explicit collation specifier:

SELECT * FROM test1 ORDER BY a || b COLLATE "fr_FR";

23.2.2. Managing Collations

A collation is an SQL schema object that maps an SQL name to locales provided by libraries installed in the operating system. A collation definition has a provider that specifies which library supplies the locale data. One standard provider name is libc, which uses the locales provided by the operating system C library. These are the locales that most tools provided by the operating system use. Another provider is ICU, which uses the external ICU library. ICU locales can only be used if support for ICU was configured when PostgreSQL was built.

A collation object provided by libc maps to a combination of LC_COLLATE and LC_CTYPE settings, as accepted by the setlocale() system library call. (As the name would suggest, the main purpose of a collation is to set LC_COLLATE, which controls the sort order. But it is rarely necessary in practice to have an LC_CTYPE setting that is different from LC_COLLATE, so it is more convenient to collect these under one concept than to create another infrastructure for setting LC_CTYPE per expression.) Also, a libc collation is tied to a character set encoding (see Section 23.3). The same collation name may exist for different encodings.

A collation object provided by ICU maps to a named collator provided by the ICU library. ICU does not support separate “collate” and “ctype” settings, so they are always the same. Also, ICU collations are independent of the encoding, so there is always only one ICU collation of a given name in a database.

23.2.2.1. Standard Collations

On all platforms, the collations named default, C, and POSIX are available. Additional collations may be available depending on operating system support. The default collation selects the LC_COLLATE and LC_CTYPE values specified at database creation time. The C and POSIX collations both specify “traditional C” behavior, in which only the ASCII letters “A” through “Z” are treated as letters, and sorting is done strictly by character code byte values.

Additionally, the SQL standard collation name ucs_basic is available for encoding UTF8. It is equivalent to C and sorts by Unicode code point.

23.2.2.2. Predefined Collations

If the operating system provides support for using multiple locales within a single program (newlocale and related functions), or if support for ICU is configured, then when a database cluster is initialized, initdb populates the system catalog pg_collation with collations based on all the locales it finds in the operating system at the time.
To inspect the currently available locales, use the query `SELECT * FROM pg_collation`, or the command `\dO+$` in psql.

### 23.2.2.2.1. libc collations

For example, the operating system might provide a locale named `de_DE.utf8`. `initdb` would then create a collation named `de_DE.utf8` for encoding UTF8 that has both `LC_COLLATE` and `LC_CTYPE` set to `de_DE.utf8`. It will also create a collation with the `.utf8` tag stripped off the name. So you could also use the collation under the name `de_DE`, which is less cumbersome to write and makes the name less encoding-dependent. Note that, nevertheless, the initial set of collation names is platform-dependent.

The default set of collations provided by `libc` map directly to the locales installed in the operating system, which can be listed using the command `locale -a`. In case a `libc` collation is needed that has different values for `LC_COLLATE` and `LC_CTYPE`, or if new locales are installed in the operating system after the database system was initialized, then a new collation may be created using the `CREATE COLLATION` command. New operating system locales can also be imported en masse using the `pg_import_system_collations()` function.

Within any particular database, only collations that use that database's encoding are of interest. Other entries in `pg_collation` are ignored. Thus, a stripped collation name such as `de_DE` can be considered unique within a given database even though it would not be unique globally. Use of the stripped collation names is recommended, since it will make one less thing you need to change if you decide to change to another database encoding. Note however that the default, `C`, and POSIX collations can be used regardless of the database encoding.

PostgreSQL considers distinct collation objects to be incompatible even when they have identical properties. Thus for example,

```
SELECT a COLLATE "C" < b COLLATE "POSIX" FROM test1;
```

will draw an error even though the `C` and `POSIX` collations have identical behaviors. Mixing stripped and non-stripped collation names is therefore not recommended.

### 23.2.2.2. ICU collations

With ICU, it is not sensible to enumerate all possible locale names. ICU uses a particular naming system for locales, but there are many more ways to name a locale than there are actually distinct locales. `initdb` uses the ICU APIs to extract a set of distinct locales to populate the initial set of collations. Collations provided by ICU are created in the SQL environment with names in BCP 47 language tag format, with a “private use” extension `-x-icu` appended, to distinguish them from `libc` locales.

Here are some example collations that might be created:

- **de-x-icu**
  
  German collation, default variant

- **de-AT-x-icu**
  
  German collation for Austria, default variant

  (There are also, say, `de-DE-x-icu` or `de-CH-x-icu`, but as of this writing, they are equivalent to `de-x-icu`.)

- **und-x-icu** (for “undefined”)

  ICU “root” collation. Use this to get a reasonable language-agnostic sort order.

---

632
Some (less frequently used) encodings are not supported by ICU. When the database encoding is one of these, ICU collation entries in `pg_collation` are ignored. Attempting to use one will draw an error along the lines of “collation "de-x-icu" for encoding "WIN874" does not exist”.

### 23.2.2.3. Creating New Collation Objects

If the standard and predefined collations are not sufficient, users can create their own collation objects using the SQL command `CREATE COLLATION`.

The standard and predefined collations are in the schema `pg_catalog`, like all predefined objects. User-defined collations should be created in user schemas. This also ensures that they are saved by `pg_dump`.

#### 23.2.2.3.1. libc collations

New libc collations can be created like this:

```sql
CREATE COLLATION german (provider = libc, locale = 'de_DE');
```

The exact values that are acceptable for the `locale` clause in this command depend on the operating system. On Unix-like systems, the command `locale -a` will show a list.

Since the predefined libc collations already include all collations defined in the operating system when the database instance is initialized, it is not often necessary to manually create new ones. Reasons might be if a different naming system is desired (in which case see also Section 23.2.2.3.3) or if the operating system has been upgraded to provide new locale definitions (in which case see also `pg_import_system_collations()`).

#### 23.2.2.3.2. ICU collations

ICU allows collations to be customized beyond the basic language+country set that is preloaded by `initdb`. Users are encouraged to define their own collation objects that make use of these facilities to suit the sorting behavior to their requirements. See [http://userguide.icu-project.org/locale](http://userguide.icu-project.org/locale) and [http://userguide.icu-project.org/collation/api](http://userguide.icu-project.org/collation/api) for information on ICU locale naming. The set of acceptable names and attributes depends on the particular ICU version.

Here are some examples:

```sql
CREATE COLLATION "de-u-co-phonebk-x-icu" (provider = icu, locale = 'de-u-co-phonebk');
CREATE COLLATION "de-u-co-phonebk-x-icu" (provider = icu, locale = '@collation=phonebook');
```

German collation with phone book collation type

The first example selects the ICU locale using a “language tag” per BCP 47. The second example uses the traditional ICU-specific locale syntax. The first style is preferred going forward, but it is not supported by older ICU versions.

Note that you can name the collation objects in the SQL environment anything you want. In this example, we follow the naming style that the predefined collations use, which in turn also follow BCP 47, but that is not required for user-defined collations.

```sql
CREATE COLLATION "und-u-co-emoji-x-icu" (provider = icu, locale = 'und-u-co-emoji');
CREATE COLLATION "und-u-co-emoji-x-icu" (provider = icu, locale = '@collation=emoji');
```

Root collation with Emoji collation type, per Unicode Technical Standard #51

Observe how in the traditional ICU locale naming system, the root locale is selected by an empty string.
CREATE COLLATION digitslast (provider = icu, locale = 'en-u-kr-latn-digit');
CREATE COLLATION digitslast (provider = icu, locale = 'en@colRe-order=latn-digit');

Sort digits after Latin letters. (The default is digits before letters.)

CREATE COLLATION upperfirst (provider = icu, locale = 'en-u-kf-upper');
CREATE COLLATION upperfirst (provider = icu, locale = 'en@colCase-First=upper');

Sort upper-case letters before lower-case letters. (The default is lower-case letters first.)

CREATE COLLATION special (provider = icu, locale = 'en-u-kr-upper-latn-digit');
CREATE COLLATION special (provider = icu, locale = 'en@colCase-First=upper;colReorder=latn-digit');

Combines both of the above options.

CREATE COLLATION numeric (provider = icu, locale = 'en-u-kn-true');
CREATE COLLATION numeric (provider = icu, locale = 'en@colNumeric=yes');

Numeric ordering, sorts sequences of digits by their numeric value, for example: A-21 < A-123
(also known as natural sort).

See Unicode Technical Standard #35\(^1\) and BCP 47\(^2\) for details. The list of possible collation types (co subtag) can be found in the CLDR repository\(^3\). The ICU Locale Explorer\(^4\) can be used to check the details of a particular locale definition. The examples using the k* subtags require at least ICU version 54.

Note that while this system allows creating collations that “ignore case” or “ignore accents” or similar (using the ks key), PostgreSQL does not at the moment allow such collations to act in a truly case- or accent-insensitive manner. Any strings that compare equal according to the collation but are not byte-wise equal will be sorted according to their byte values.

By design, ICU will accept almost any string as a locale name and match it to the closest locale it can provide, using the fallback procedure described in its documentation. Thus, there will be no direct feedback if a collation specification is composed using features that the given ICU installation does not actually support. It is therefore recommended to create application-level test cases to check that the collation definitions satisfy one's requirements.

### 23.2.2.3.3. Copying Collations

The command CREATE COLLATION can also be used to create a new collation from an existing collation, which can be useful to be able to use operating-system-independent collation names in applications, create compatibility names, or use an ICU-provided collation under a more readable name. For example:

CREATE COLLATION german FROM "de_DE";

---

\(^1\) [http://unicode.org/reports/tr35/tr35-collation.html](http://unicode.org/reports/tr35/tr35-collation.html)


\(^3\) [http://www.unicode.org/repos/cldr/trunk/common/bcp47/collation.xml](http://www.unicode.org/repos/cldr/trunk/common/bcp47/collation.xml)

\(^4\) [https://ssl.icu-project.org/icu-bin/locexp](https://ssl.icu-project.org/icu-bin/locexp)
CREATE COLLATION french FROM "fr-x-icu";

# 23.3. Character Set Support

The character set support in PostgreSQL allows you to store text in a variety of character sets (also called encodings), including single-byte character sets such as the ISO 8859 series and multiple-byte character sets such as EUC (Extended Unix Code), UTF-8, and Mule internal code. All supported character sets can be used transparently by clients, but a few are not supported for use within the server (that is, as a server-side encoding). The default character set is selected while initializing your PostgreSQL database cluster using `initdb`. It can be overridden when you create a database, so you can have multiple databases each with a different character set.

An important restriction, however, is that each database's character set must be compatible with the database's `LC_CTYPE` (character classification) and `LC_COLLATE` (string sort order) locale settings. For C or POSIX locale, any character set is allowed, but for other libc-provided locales there is only one character set that will work correctly. (On Windows, however, UTF-8 encoding can be used with any locale.) If you have ICU support configured, ICU-provided locales can be used with most but not all server-side encodings.

## 23.3.1. Supported Character Sets

Table 23.1 shows the character sets available for use in PostgreSQL.

### Table 23.1. PostgreSQL Character Sets

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Language</th>
<th>Server?</th>
<th>ICU?</th>
<th>Bytes/Char</th>
<th>Aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>BIG5</td>
<td>Big Five</td>
<td>Traditional Chinese</td>
<td>No</td>
<td>No</td>
<td>1-2</td>
<td>WIN950, Windows950</td>
</tr>
<tr>
<td>EUC_CN</td>
<td>Extended UNIX Code-CN</td>
<td>Simplified Chinese</td>
<td>Yes</td>
<td>Yes</td>
<td>1-3</td>
<td></td>
</tr>
<tr>
<td>EU-C_JIS_200</td>
<td>Extended UNIX Code-JP</td>
<td>Japanese</td>
<td>Yes</td>
<td>No</td>
<td>1-3</td>
<td></td>
</tr>
<tr>
<td>EUC_KR</td>
<td>Extended UNIX Code-KR</td>
<td>Korean</td>
<td>Yes</td>
<td>Yes</td>
<td>1-3</td>
<td></td>
</tr>
<tr>
<td>EUC_TW</td>
<td>Extended UNIX Code-TW</td>
<td>Traditional Chinese, Taiwanese</td>
<td>Yes</td>
<td>Yes</td>
<td>1-3</td>
<td></td>
</tr>
<tr>
<td>GB18030</td>
<td>National Standard</td>
<td>Chinese</td>
<td>No</td>
<td>No</td>
<td>1-4</td>
<td></td>
</tr>
<tr>
<td>GBK</td>
<td>Extended National Standard</td>
<td>Simplified Chinese</td>
<td>No</td>
<td>No</td>
<td>1-2</td>
<td>WIN936, Windows936</td>
</tr>
<tr>
<td>ISO_8859_5</td>
<td>ISO 8859-5, Latin/Cyrillic</td>
<td>Yes</td>
<td>Yes</td>
<td></td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>ISO_8859_6</td>
<td>ISO 8859-6, Latin/Arabic</td>
<td>Yes</td>
<td>Yes</td>
<td></td>
<td>1</td>
<td></td>
</tr>
</tbody>
</table>

635
<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Language</th>
<th>Server?</th>
<th>ICU?</th>
<th>Bytes/Char</th>
<th>Aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>ISO_8859-7</td>
<td>ISO 8859-7, Latin/Greek</td>
<td></td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td></td>
<td>ECMA 118</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ISO_8859-8</td>
<td>ISO 8859-8, Latin/</td>
<td></td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Hebrew ECMA 121</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>JOHAB</td>
<td>JOHAB</td>
<td>Korean (Hangul)</td>
<td>No</td>
<td>No</td>
<td>1-3</td>
<td></td>
</tr>
<tr>
<td>KOI8R</td>
<td>KOI8-R</td>
<td>Cyrillic (Russian)</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td>KOI8</td>
</tr>
<tr>
<td>KOI8U</td>
<td>KOI8-U</td>
<td>Cyrillic (Ukrainian)</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>LATIN1</td>
<td>ISO 8859-1, Western European</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO88591</td>
</tr>
<tr>
<td></td>
<td>ECMA 94</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN2</td>
<td>ISO 8859-2, Central European</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO88592</td>
</tr>
<tr>
<td></td>
<td>ECMA 94</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN3</td>
<td>ISO 8859-3, South European</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO88593</td>
</tr>
<tr>
<td></td>
<td>ECMA 94</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN4</td>
<td>ISO 8859-4, North European</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO88594</td>
</tr>
<tr>
<td></td>
<td>ECMA 94</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN5</td>
<td>ISO 8859-9, Turkish</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO88599</td>
</tr>
<tr>
<td></td>
<td>ECMA 128</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN6</td>
<td>ISO 8859-10, Nordic</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO885910</td>
</tr>
<tr>
<td></td>
<td>ECMA 144</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LATIN7</td>
<td>ISO 8859-13, Baltic</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO885913</td>
</tr>
<tr>
<td>LATIN8</td>
<td>ISO 8859-14, Celtic</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO885914</td>
</tr>
<tr>
<td>LATIN9</td>
<td>ISO 8859-15, LATIN1 with Euro and accents</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1</td>
<td>ISO885915</td>
</tr>
<tr>
<td>LATIN10</td>
<td>ISO 8859-16, Romanian</td>
<td></td>
<td>Yes</td>
<td>No</td>
<td>1</td>
<td>ISO885916</td>
</tr>
<tr>
<td></td>
<td>ASRO SR 14111</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>MULE_INTERNAL</td>
<td>Mule internal code</td>
<td>Multilingual</td>
<td>Yes</td>
<td>No</td>
<td>1-4</td>
<td></td>
</tr>
<tr>
<td>SJIS</td>
<td>Shift JIS</td>
<td>Japanese</td>
<td>No</td>
<td>No</td>
<td>1-2</td>
<td>Mskanji, ShiftJIS, WIN932, Windows932</td>
</tr>
<tr>
<td>SHIFT_JIS</td>
<td>Shift JIS</td>
<td>Japanese</td>
<td>No</td>
<td>No</td>
<td>1-2</td>
<td></td>
</tr>
<tr>
<td>SQL_ASCII</td>
<td>unspecified (see text)</td>
<td>any</td>
<td>Yes</td>
<td>No</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>UHC</td>
<td>Unified Hangul Code</td>
<td>Korean</td>
<td>No</td>
<td>No</td>
<td>1-2</td>
<td>WIN949, Windows949</td>
</tr>
<tr>
<td>UTF8</td>
<td>Unicode, 8-bit</td>
<td></td>
<td>Yes</td>
<td></td>
<td>1-4</td>
<td>Unicode</td>
</tr>
</tbody>
</table>
## Localization

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Language</th>
<th>Server?</th>
<th>ICU?</th>
<th>Bytes/Char</th>
<th>Aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>WIN866</td>
<td>Windows CP866</td>
<td>Cyrillic</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td>ALT</td>
</tr>
<tr>
<td>WIN874</td>
<td>Windows CP864</td>
<td>Thai</td>
<td>Yes</td>
<td>No</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1250</td>
<td>Windows CP1250</td>
<td>Central European</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td>WIN</td>
</tr>
<tr>
<td>WIN1251</td>
<td>Windows CP1251</td>
<td>Cyrillic</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td>WIN</td>
</tr>
<tr>
<td>WIN1252</td>
<td>Windows CP1252</td>
<td>Western European</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1253</td>
<td>Windows CP1253</td>
<td>Greek</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1254</td>
<td>Windows CP1254</td>
<td>Turkish</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1255</td>
<td>Windows CP1255</td>
<td>Hebrew</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1256</td>
<td>Windows CP1256</td>
<td>Arabic</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1257</td>
<td>Windows CP1257</td>
<td>Baltic</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>WIN1258</td>
<td>Windows CP1258</td>
<td>Vietnamese</td>
<td>Yes</td>
<td>Yes</td>
<td>1</td>
<td>ABC, TCVN, TCVN5712, VSCII</td>
</tr>
</tbody>
</table>

Not all client APIs support all the listed character sets. For example, the PostgreSQL JDBC driver does not support MULE INTERNAL, LATIN6, LATIN8, and LATIN10.

The `SQL_ASCII` setting behaves considerably differently from the other settings. When the server character set is `SQL_ASCII`, the server interprets byte values 0-127 according to the ASCII standard, while byte values 128-255 are taken as uninterpreted characters. No encoding conversion will be done when the setting is `SQL_ASCII`. Thus, this setting is not so much a declaration that a specific encoding is in use, as a declaration of ignorance about the encoding. In most cases, if you are working with any non-ASCII data, it is unwise to use the `SQL_ASCII` setting because PostgreSQL will be unable to help you by converting or validating non-ASCII characters.

### 23.3.2. Setting the Character Set

`initdb` defines the default character set (encoding) for a PostgreSQL cluster. For example,

```
initdb -E EUC_JP
```

sets the default character set to EUC_JP (Extended Unix Code for Japanese). You can use `--encoding` instead of `-E` if you prefer longer option strings. If no `-E` or `--encoding` option is given, `initdb` attempts to determine the appropriate encoding to use based on the specified or default locale.

You can specify a non-default encoding at database creation time, provided that the encoding is compatible with the selected locale:

```
createdb -E EUC_KR -T template0 --lc-collate=ko_KR.euckr --lc-ctype=ko_KR.euckr korean
```
This will create a database named `korean` that uses the character set `EUC_KR`, and locale `ko_KR`. Another way to accomplish this is to use this SQL command:

```
CREATE DATABASE korean WITH ENCODING 'EUC_KR'
  LC_COLLATE='ko_KR.euckr' LC_CTYPE='ko_KR.euckr'
  TEMPLATE=template0;
```

Notice that the above commands specify copying the `template0` database. When copying any other database, the encoding and locale settings cannot be changed from those of the source database, because that might result in corrupt data. For more information see Section 22.3.

The encoding for a database is stored in the system catalog `pg_database`. You can see it by using the `psql -l` option or the `\l` command.

```
$ psql -l
```

<table>
<thead>
<tr>
<th>Name</th>
<th>Owner</th>
<th>Encoding</th>
<th>Collation</th>
<th>Ctype</th>
</tr>
</thead>
<tbody>
<tr>
<td>clocaledb</td>
<td>hlinnaka</td>
<td>SQL_ASCII</td>
<td>C</td>
<td>C</td>
</tr>
<tr>
<td>englishdb</td>
<td>hlinnaka</td>
<td>UTF8</td>
<td>en_GB.UTF8</td>
<td>en_GB.UTF8</td>
</tr>
<tr>
<td>japanese</td>
<td>hlinnaka</td>
<td>UTF8</td>
<td>ja_JP.UTF8</td>
<td>ja_JP.UTF8</td>
</tr>
<tr>
<td>korean</td>
<td>hlinnaka</td>
<td>EUC_KR</td>
<td>ko_KR.euckr</td>
<td>ko_KR.euckr</td>
</tr>
<tr>
<td>postgres</td>
<td>hlinnaka</td>
<td>UTF8</td>
<td>fi_FI.UTF8</td>
<td>fi_FI.UTF8</td>
</tr>
<tr>
<td>template0</td>
<td>hlinnaka</td>
<td>UTF8</td>
<td>fi_FI.UTF8</td>
<td>fi_FI.UTF8</td>
</tr>
</tbody>
</table>

(7 rows)

### Important

On most modern operating systems, PostgreSQL can determine which character set is implied by the `LC_CTYPE` setting, and it will enforce that only the matching database encoding is used. On older systems it is your responsibility to ensure that you use the encoding expected by the locale you have selected. A mistake in this area is likely to lead to strange behavior of locale-dependent operations such as sorting.

PostgreSQL will allow superusers to create databases with `SQL_ASCII` encoding even when `LC_CTYPE` is not `C` or `POSIX`. As noted above, `SQL_ASCII` does not enforce that the data stored in the database has any particular encoding, and so this choice poses risks of locale-dependent misbehavior. Using this combination of settings is deprecated and may someday be forbidden altogether.

### 23.3.3. Automatic Character Set Conversion Between Server and Client

PostgreSQL supports automatic character set conversion between server and client for certain character set combinations. The conversion information is stored in the `pg_conversion` system catalog. PostgreSQL comes with some predefined conversions, as shown in Table 23.2. You can create a new conversion using the SQL command `CREATE CONVERSION`.

638
### Table 23.2. Client/Server Character Set Conversions

<table>
<thead>
<tr>
<th>Server Character Set</th>
<th>Available Client Character Sets</th>
</tr>
</thead>
<tbody>
<tr>
<td>BIG5</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>EUC_CN</td>
<td>EUC_CN, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>EUC_JP</td>
<td>EUC_JP, MULE_INTERNAL, SJIS, UTF8</td>
</tr>
<tr>
<td>EUC_JIS_2004</td>
<td>EUC_JIS_2004, SHIFT_JIS_2004, UTF8</td>
</tr>
<tr>
<td>EUC_KR</td>
<td>EUC_KR, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>EUC_TW</td>
<td>EUC_TW, BIG5, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>GB18030</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>GBK</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>ISO_8859_5</td>
<td>ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN866, WIN1251</td>
</tr>
<tr>
<td>ISO_8859_6</td>
<td>ISO_8859_6, UTF8</td>
</tr>
<tr>
<td>ISO_8859_7</td>
<td>ISO_8859_7, UTF8</td>
</tr>
<tr>
<td>ISO_8859_8</td>
<td>ISO_8859_8, UTF8</td>
</tr>
<tr>
<td>JOHAB</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>KOI8R</td>
<td>KOI8R, ISO_8859_5, MULE_INTERNAL, UTF8, WIN866, WIN1251</td>
</tr>
<tr>
<td>KOI8U</td>
<td>KOI8U, UTF8</td>
</tr>
<tr>
<td>LATIN1</td>
<td>LATIN1, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN2</td>
<td>LATIN2, MULE_INTERNAL, UTF8, WIN1250</td>
</tr>
<tr>
<td>LATIN3</td>
<td>LATIN3, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN4</td>
<td>LATIN4, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN5</td>
<td>LATIN5, UTF8</td>
</tr>
<tr>
<td>LATIN6</td>
<td>LATIN6, UTF8</td>
</tr>
<tr>
<td>LATIN7</td>
<td>LATIN7, UTF8</td>
</tr>
<tr>
<td>LATIN8</td>
<td>LATIN8, UTF8</td>
</tr>
<tr>
<td>LATIN9</td>
<td>LATIN9, UTF8</td>
</tr>
<tr>
<td>LATIN10</td>
<td>LATIN10, UTF8</td>
</tr>
<tr>
<td>MULE_INTERNAL</td>
<td>MULE_INTERNAL, BIG5, EUC_CN, EUC_JP, EUC_KR, EUC_TW, ISO_8859_5, KOI8R, LATIN1 to LATIN4, SJIS, WIN866, WIN1250, WIN1251</td>
</tr>
<tr>
<td>SJIS</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>SHIFT_JIS_2004</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>SQL_ASCII</td>
<td>any (no conversion will be performed)</td>
</tr>
<tr>
<td>UTF8</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>WIN866</td>
<td>WIN866, ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN1251</td>
</tr>
<tr>
<td>WIN874</td>
<td>WIN874, UTF8</td>
</tr>
<tr>
<td>WIN1250</td>
<td>WIN1250, LATIN2, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>WIN1251</td>
<td>WIN1251, ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN866</td>
</tr>
</tbody>
</table>
To enable automatic character set conversion, you have to tell PostgreSQL the character set (encoding) you would like to use in the client. There are several ways to accomplish this:

- Using the `\encoding` command in psql. `\encoding` allows you to change client encoding on the fly. For example, to change the encoding to SJIS, type:

  \encoding SJIS

- libpq (Section 34.10) has functions to control the client encoding.

- Using `SET client_encoding TO`. Setting the client encoding can be done with this SQL command:

  SET CLIENT_ENCODING TO 'value';

  Also you can use the standard SQL syntax `SET NAMES` for this purpose:

  SET NAMES 'value';

  To query the current client encoding:

  SHOW client_encoding;

  To return to the default encoding:

  RESET client_encoding;

- Using `PGCLIENTENCODING`. If the environment variable `PGCLIENTENCODING` is defined in the client's environment, that client encoding is automatically selected when a connection to the server is made. (This can subsequently be overridden using any of the other methods mentioned above.)

- Using the configuration variable `client_encoding`. If the `client_encoding` variable is set, that client encoding is automatically selected when a connection to the server is made. (This can subsequently be overridden using any of the other methods mentioned above.)

If the conversion of a particular character is not possible — suppose you chose EUC_JP for the server and LATIN1 for the client, and some Japanese characters are returned that do not have a representation in LATIN1 — an error is reported.

If the client character set is defined as SQL_ASCII, encoding conversion is disabled, regardless of the server's character set. Just as for the server, use of SQL_ASCII is unwise unless you are working with all-ASCII data.
23.3.4. Further Reading

These are good sources to start learning about various kinds of encoding systems.

*CJKV Information Processing: Chinese, Japanese, Korean & Vietnamese Computing*

Contains detailed explanations of EUC_JP, EUC_CN, EUC_KR, EUC_TW.

http://www.unicode.org/

The web site of the Unicode Consortium.

RFC 3629

UTF-8 (8-bit UCS/Unicode Transformation Format) is defined here.
Chapter 24. Routine Database Maintenance Tasks

PostgreSQL, like any database software, requires that certain tasks be performed regularly to achieve optimum performance. The tasks discussed here are required, but they are repetitive in nature and can easily be automated using standard tools such as cron scripts or Windows' Task Scheduler. It is the database administrator's responsibility to set up appropriate scripts, and to check that they execute successfully.

One obvious maintenance task is the creation of backup copies of the data on a regular schedule. Without a recent backup, you have no chance of recovery after a catastrophe (disk failure, fire, mistakenly dropping a critical table, etc.). The backup and recovery mechanisms available in PostgreSQL are discussed at length in Chapter 25.

The other main category of maintenance task is periodic “vacuuming” of the database. This activity is discussed in Section 24.1. Closely related to this is updating the statistics that will be used by the query planner, as discussed in Section 24.1.3.

Another task that might need periodic attention is log file management. This is discussed in Section 24.3.

check_postgres is available for monitoring database health and reporting unusual conditions. check_postgres integrates with Nagios and MRTG, but can be run standalone too.

PostgreSQL is low-maintenance compared to some other database management systems. Nonetheless, appropriate attention to these tasks will go far towards ensuring a pleasant and productive experience with the system.

24.1. Routine Vacuuming

PostgreSQL databases require periodic maintenance known as vacuuming. For many installations, it is sufficient to let vacuuming be performed by the autovacuum daemon, which is described in Section 24.1.6. You might need to adjust the autovacuuming parameters described there to obtain best results for your situation. Some database administrators will want to supplement or replace the daemon's activities with manually-managed VACUUM commands, which typically are executed according to a schedule by cron or Task Scheduler scripts. To set up manually-managed vacuuming properly, it is essential to understand the issues discussed in the next few subsections. Administrators who rely on autovacuuming may still wish to skim this material to help them understand and adjust autovacuuming.

24.1.1. Vacuuming Basics

PostgreSQL’s VACUUM command has to process each table on a regular basis for several reasons:

1. To recover or reuse disk space occupied by updated or deleted rows.
2. To update data statistics used by the PostgreSQL query planner.
3. To update the visibility map, which speeds up index-only scans.
4. To protect against loss of very old data due to transaction ID wraparound or multixact ID wraparound.

Each of these reasons dictates performing VACUUM operations of varying frequency and scope, as explained in the following subsections.

---

1 https://bucardo.org/check_postgres/
There are two variants of VACUUM: standard VACUUM and VACUUM FULL. VACUUM FULL can reclaim more disk space but runs much more slowly. Also, the standard form of VACUUM can run in parallel with production database operations. (Commands such as SELECT, INSERT, UPDATE, and DELETE will continue to function normally, though you will not be able to modify the definition of a table with commands such as ALTER TABLE while it is being vacuumed.) VACUUM FULL requires exclusive lock on the table it is working on, and therefore cannot be done in parallel with other use of the table. Generally, therefore, administrators should strive to use standard VACUUM and avoid VACUUM FULL.

VACUUM creates a substantial amount of I/O traffic, which can cause poor performance for other active sessions. There are configuration parameters that can be adjusted to reduce the performance impact of background vacuuming — see Section 19.4.4.

### 24.1.2. Recovering Disk Space

In PostgreSQL, an UPDATE or DELETE of a row does not immediately remove the old version of the row. This approach is necessary to gain the benefits of multiversion concurrency control (MVCC, see Chapter 13): the row version must not be deleted while it is still potentially visible to other transactions. But eventually, an outdated or deleted row version is no longer of interest to any transaction. The space it occupies must then be reclaimed for reuse by new rows, to avoid unbounded growth of disk space requirements. This is done by running VACUUM.

The standard form of VACUUM removes dead row versions in tables and indexes and marks the space available for future reuse. However, it will not return the space to the operating system, except in the special case where one or more pages at the end of a table become entirely free and an exclusive table lock can be easily obtained. In contrast, VACUUM FULL actively compacts tables by writing a complete new version of the table file with no dead space. This minimizes the size of the table, but can take a long time. It also requires extra disk space for the new copy of the table, until the operation completes.

The usual goal of routine vacuuming is to do standard VACUUMs often enough to avoid needing VACUUM FULL. The autovacuum daemon attempts to work this way, and in fact will never issue VACUUM FULL. In this approach, the idea is not to keep tables at their minimum size, but to maintain steady-state usage of disk space: each table occupies space equivalent to its minimum size plus however much space gets used up between vacuumings. Although VACUUM FULL can be used to shrink a table back to its minimum size and return the disk space to the operating system, there is not much point in this if the table will just grow again in the future. Thus, moderately-frequent standard VACUUM runs are a better approach than infrequent VACUUM FULL runs for maintaining heavily-updated tables.

Some administrators prefer to schedule vacuuming themselves, for example doing all the work at night when load is low. The difficulty with doing vacuuming according to a fixed schedule is that if a table has an unexpected spike in update activity, it may get bloated to the point that VACUUM FULL is really necessary to reclaim space. Using the autovacuum daemon alleviates this problem, since the daemon schedules vacuuming dynamically in response to update activity. It is unwise to disable the daemon completely unless you have an extremely predictable workload. One possible compromise is to set the daemon's parameters so that it will only react to unusually heavy update activity, thus keeping things from getting out of hand, while scheduled VACUUMs are expected to do the bulk of the work when the load is typical.

For those not using autovacuum, a typical approach is to schedule a database-wide VACUUM once a day during a low-usage period, supplemented by more frequent vacuuming of heavily-updated tables as necessary. (Some installations with extremely high update rates vacuum their busiest tables as often as once every few minutes.) If you have multiple databases in a cluster, don't forget to VACUUM each one; the program vacuumdb might be helpful.

**Tip**

Plain VACUUM may not be satisfactory when a table contains large numbers of dead row versions as a result of massive update or delete activity. If you have such a table...
and you need to reclaim the excess disk space it occupies, you will need to use VACUUM FULL, or alternatively CLUSTER or one of the table-rewriting variants of ALTER TABLE. These commands rewrite an entire new copy of the table and build new indexes for it. All these options require exclusive lock. Note that they also temporarily use extra disk space approximately equal to the size of the table, since the old copies of the table and indexes can’t be released until the new ones are complete.

**Tip**

If you have a table whose entire contents are deleted on a periodic basis, consider doing it with TRUNCATE rather than using DELETE followed by VACUUM. TRUNCATE removes the entire content of the table immediately, without requiring a subsequent VACUUM or VACUUM FULL to reclaim the now-unused disk space. The disadvantage is that strict MVCC semantics are violated.

### 24.1.3. Updating Planner Statistics

The PostgreSQL query planner relies on statistical information about the contents of tables in order to generate good plans for queries. These statistics are gathered by the ANALYZE command, which can be invoked by itself or as an optional step in VACUUM. It is important to have reasonably accurate statistics, otherwise poor choices of plans might degrade database performance.

The autovacuum daemon, if enabled, will automatically issue ANALYZE commands whenever the content of a table has changed sufficiently. However, administrators might prefer to rely on manually-scheduled ANALYZE operations, particularly if it is known that update activity on a table will not affect the statistics of “interesting” columns. The daemon schedules ANALYZE strictly as a function of the number of rows inserted or updated; it has no knowledge of whether that will lead to meaningful statistical changes.

As with vacuuming for space recovery, frequent updates of statistics are more useful for heavily-updated tables than for seldom-updated ones. But even for a heavily-updated table, there might be no need for statistics updates if the statistical distribution of the data is not changing much. A simple rule of thumb is to think about how much the minimum and maximum values of the columns in the table change. For example, a timestamp column that contains the time of row update will have a constantly-increasing maximum value as rows are added and updated; such a column will probably need more frequent statistics updates than, say, a column containing URLs for pages accessed on a website. The URL column might receive changes just as often, but the statistical distribution of its values probably changes relatively slowly.

It is possible to run ANALYZE on specific tables and even just specific columns of a table, so the flexibility exists to update some statistics more frequently than others if your application requires it. In practice, however, it is usually best to just analyze the entire database, because it is a fast operation. ANALYZE uses a statistically random sampling of the rows of a table rather than reading every single row.

**Tip**

Although per-column tweaking of ANALYZE frequency might not be very productive, you might find it worthwhile to do per-column adjustment of the level of detail of the statistics collected by ANALYZE. Columns that are heavily used in WHERE clauses and have highly irregular data distributions might require a finer-grain data histogram than other columns. See ALTER TABLE SET STATISTICS, or change the database-wide default using the default_statistics_target configuration parameter.

Also, by default there is limited information available about the selectivity of functions. However, if you create an expression index that uses a function call, useful statistics
will be gathered about the function, which can greatly improve query plans that use the expression index.

**Tip**

The autovacuum daemon does not issue `ANALYZE` commands for foreign tables, since it has no means of determining how often that might be useful. If your queries require statistics on foreign tables for proper planning, it's a good idea to run manually-managed `ANALYZE` commands on those tables on a suitable schedule.

### 24.1.4. Updating The Visibility Map

Vacuum maintains a visibility map for each table to keep track of which pages contain only tuples that are known to be visible to all active transactions (and all future transactions, until the page is again modified). This has two purposes. First, vacuum itself can skip such pages on the next run, since there is nothing to clean up.

Second, it allows PostgreSQL to answer some queries using only the index, without reference to the underlying table. Since PostgreSQL indexes don't contain tuple visibility information, a normal index scan fetches the heap tuple for each matching index entry, to check whether it should be seen by the current transaction. An *index-only scan*, on the other hand, checks the visibility map first. If it's known that all tuples on the page are visible, the heap fetch can be skipped. This is most useful on large data sets where the visibility map can prevent disk accesses. The visibility map is vastly smaller than the heap, so it can easily be cached even when the heap is very large.

### 24.1.5. Preventing Transaction ID Wraparound Failures

PostgreSQL’s MVCC transaction semantics depend on being able to compare transaction ID (XID) numbers: a row version with an insertion XID greater than the current transaction’s XID is “in the future” and should not be visible to the current transaction. But since transaction IDs have limited size (32 bits) a cluster that runs for a long time (more than 4 billion transactions) would suffer *transaction ID wraparound*: the XID counter wraps around to zero, and all of a sudden transactions that were in the past appear to be in the future — which means their output become invisible. In short, catastrophic data loss. (Actually the data is still there, but that’s cold comfort if you cannot get at it.) To avoid this, it is necessary to vacuum every table in every database at least once every two billion transactions.

The reason that periodic vacuuming solves the problem is that `VACUUM` will mark rows as *frozen*, indicating that they were inserted by a transaction that committed sufficiently far in the past that the effects of the inserting transaction are certain to be visible to all current and future transactions. Normal XIDs are compared using modulo-2\(^{32}\) arithmetic. This means that for every normal XID, there are two billion XIDs that are “older” and two billion that are “newer”; another way to say it is that the normal XID space is circular with no endpoint. Therefore, once a row version has been created with a particular normal XID, the row version will appear to be “in the past” for the next two billion transactions, no matter which normal XID we are talking about. If the row version still exists after more than two billion transactions, it will suddenly appear to be in the future. To prevent this, PostgreSQL reserves a special XID, *FrozenTransactionId*, which does not follow the normal XID comparison rules and is always considered older than every normal XID. Frozen row versions are treated as if the inserting XID were *FrozenTransactionId*, so that they will appear to be “in the past” to all normal transactions regardless of wraparound issues, and so such row versions will be valid until deleted, no matter how long that is.

**Note**

In PostgreSQL versions before 9.4, freezing was implemented by actually replacing a row’s insertion XID with *FrozenTransactionId*, which was visible in the row’s
xmin system column. Newer versions just set a flag bit, preserving the row's original xmin for possible forensic use. However, rows with xmin equal to FrozenTransactionId (2) may still be found in databases pg_upgrade'd from pre-9.4 versions.

Also, system catalogs may contain rows with xmin equal to BootstrapTransactionId (1), indicating that they were inserted during the first phase of initdb. Like FrozenTransactionId, this special XID is treated as older than every normal XID.

vacuum_freeze_min_age controls how old an XID value has to be before rows bearing that XID will be frozen. Increasing this setting may avoid unnecessary work if the rows that would otherwise be frozen will soon be modified again, but decreasing this setting increases the number of transactions that can elapse before the table must be vacuumed again.

VACUUM uses the visibility map to determine which pages of a table must be scanned. Normally, it will skip pages that don't have any dead row versions even if those pages might still have row versions with old XID values. Therefore, normal VACUUMs won't always freeze every old row version in the table. Periodically, VACUUM will perform an aggressive vacuum, skipping only those pages which contain neither dead rows nor any unfrozen XID or MXID values. vacuum_freeze_table_age controls when VACUUM does that: all-visible but not all-frozen pages are scanned if the number of transactions that have passed since the last such scan is greater than vacuum_freeze_table_age minus vacuum_freeze_min_age. Setting vacuum_freeze_table_age to 0 forces VACUUM to use this more aggressive strategy for all scans.

The maximum time that a table can go unvacuumed is two billion transactions minus the vacuum_freeze_min_age value at the time of the last aggressive vacuum. If it were to go unvacuumed for longer than that, data loss could result. To ensure that this does not happen, autovacuum is invoked on any table that might contain unfrozen rows with XIDs older than the age specified by the configuration parameter autovacuum_freeze_max_age. (This will happen even if autovacuum is disabled.)

This implies that if a table is not otherwise vacuumed, autovacuum will be invoked on it approximately once every autovacuum_freeze_max_age minus vacuum_freeze_min_age transactions. For tables that are regularly vacuumed for space reclamation purposes, this is of little importance. However, for static tables (including tables that receive inserts, but no updates or deletes), there is no need to vacuum for space reclamation, so it can be useful to try to maximize the interval between forced autovacuums on very large static tables. Obviously one can do this either by increasing autovacuum_freeze_max_age or decreasing vacuum_freeze_min_age.

The effective maximum for vacuum_freeze_table_age is 0.95 * autovacuum_freeze_max_age; a setting higher than that will be capped to the maximum. A value higher than autovacuum_freeze_max_age wouldn't make sense because an anti-wraparound autovacuum would be triggered at that point anyway, and the 0.95 multiplier leaves some breathing room to run a manual VACUUM before that happens. As a rule of thumb, vacuum_freeze_table_age should be set to a value somewhat below autovacuum_freeze_max_age, leaving enough gap so that a regularly scheduled VACUUM or an autovacuum triggered by normal delete and update activity is run in that window. Setting it too close could lead to anti-wraparound autovacuums, even though the table was recently vacuumed to reclaim space, whereas lower values lead to more frequent aggressive vacuuming.

The sole disadvantage of increasing autovacuum_freeze_max_age (and vacuum_freeze_table_age along with it) is that the pg_xact and pg_commit_ts subdirectories of the database cluster will take more space, because it must store the commit status and (if track_commit_timestamp is enabled) timestamp of all transactions back to the autovacuum_freeze_max_age horizon. The commit status uses two bits per transaction, so if autovacuum_freeze_max_age is set to its maximum allowed value of two billion, pg_xact can be expected to grow to about half a gigabyte and pg_commit_ts to about 20GB. If this is trivial compared to your total database size, setting autovacuum_freeze_max_age to its maximum allowed value is recommended. Otherwise, set it depending on what you are willing to allow for pg_xact and

646
Routine Database
Maintenance Tasks

pg_commit_ts storage. (The default, 200 million transactions, translates to about 50MB of pg_xact storage and about 2GB of pg_commit_ts storage.)

One disadvantage of decreasing vacuum_freeze_min_age is that it might cause VACUUM to do useless work: freezing a row version is a waste of time if the row is modified soon thereafter (causing it to acquire a new XID). So the setting should be large enough that rows are not frozen until they are unlikely to change any more.

To track the age of the oldest unfrozen XIDs in a database, VACUUM stores XID statistics in the system tables pg_class and pg_database. In particular, the relfrozenxid column of a table's pg_class row contains the freeze cutoff XID that was used by the last aggressive VACUUM for that table. All rows inserted by transactions with XIDs older than this cutoff XID are guaranteed to have been frozen. Similarly, the datfrozenxid column of a database's pg_database row is a lower bound on the unfrozen XIDs appearing in that database — it is just the minimum of the per-table relfrozenxid values within the database. A convenient way to examine this information is to execute queries such as:

```sql
SELECT c.oid::regclass as table_name,
       greatest(age(c.relfrozenxid),age(t.relfrozenxid)) as age
FROM pg_class c
LEFT JOIN pg_class t ON c.reltoastrelid = t.oid
WHERE c.relkind IN ('r', 'm');

SELECT datname, age(datfrozenxid) FROM pg_database;
```

The age column measures the number of transactions from the cutoff XID to the current transaction's XID.

VACUUM normally only scans pages that have been modified since the last vacuum, but relfrozenxid can only be advanced when every page of the table that might contain unfrozen XIDs is scanned. This happens when relfrozenxid is more than vacuum_freeze_table_age transactions old, when VACUUM's FREEZE option is used, or when all pages that are not already all-frozen happen to require vacuuming to remove dead row versions. When VACUUM scans every page in the table that is not already all-frozen, it should set age(relfrozenxid) to a value just a little more than the vacuum_freeze_min_age setting that was used (more by the number of transactions started since the VACUUM started). If no relfrozenxid-advancing VACUUM is issued on the table until autovacuum_freeze_max_age is reached, an autovacuum will soon be forced for the table.

If for some reason autovacuum fails to clear old XIDs from a table, the system will begin to emit warning messages like this when the database's oldest XIDs reach ten million transactions from the wraparound point:

```
WARNING:  database "mydb" must be vacuumed within 177009986 transactions
HINT:  To avoid a database shutdown, execute a database-wide VACUUM in "mydb".
```

(A manual VACUUM should fix the problem, as suggested by the hint; but note that the VACUUM must be performed by a superuser, else it will fail to process system catalogs and thus not be able to advance the database's datfrozenxid.) If these warnings are ignored, the system will shut down and refuse to start any new transactions once there are fewer than 1 million transactions left until wraparound:

```
ERROR:  database is not accepting commands to avoid wraparound data loss in database "mydb"
HINT:  Stop the postmaster and vacuum that database in single-user mode.
```
The 1-million-transaction safety margin exists to let the administrator recover without data loss, by manually executing the required VACUUM commands. However, since the system will not execute commands once it has gone into the safety shutdown mode, the only way to do this is to stop the server and start the server in single-user mode to execute VACUUM. The shutdown mode is not enforced in single-user mode. See the postgres reference page for details about using single-user mode.

24.1.5.1. Multixacts and Wraparound

Multixact IDs are used to support row locking by multiple transactions. Since there is only limited space in a tuple header to store lock information, that information is encoded as a “multiple transaction ID”, or multixact ID for short, whenever there is more than one transaction concurrently locking a row. Information about which transaction IDs are included in any particular multixact ID is stored separately in the pg_multixact subdirectory, and only the multixact ID appears in the xmax field in the tuple header. Like transaction IDs, multixact IDs are implemented as a 32-bit counter and corresponding storage, all of which requires careful aging management, storage cleanup, and wraparound handling. There is a separate storage area which holds the list of members in each multixact, which also uses a 32-bit counter and which must also be managed.

Whenever VACUUM scans any part of a table, it will replace any multixact ID it encounters which is older than vacuum_multixact_freeze_min_age by a different value, which can be the zero value, a single transaction ID, or a newer multixact ID. For each table, pg_class.relminmxid stores the oldest possible multixact ID still appearing in any tuple of that table. If this value is older than vacuum_multixact_freeze_table_age, an aggressive vacuum is forced. As discussed in the previous section, an aggressive vacuum means that only those pages which are known to be all-frozen will be skipped. mxid_age() can be used on pg_class.relminmxid to find its age.

Aggressive VACUUM scans, regardless of what causes them, enable advancing the value for that table. Eventually, as all tables in all databases are scanned and their oldest multixact values are advanced, on-disk storage for older multixacts can be removed.

As a safety device, an aggressive vacuum scan will occur for any table whose multixact-age is greater than autovacuum_multixact_freeze_max_age. Aggressive vacuum scans will also occur progressively for all tables, starting with those that have the oldest multixact-age, if the amount of used member storage space exceeds the amount 50% of the addressable storage space. Both of these kinds of aggressive scans will occur even if autovacuum is nominally disabled.

24.1.6. The Autovacuum Daemon

PostgreSQL has an optional but highly recommended feature called autovacuum, whose purpose is to automate the execution of VACUUM and ANALYZE commands. When enabled, autovacuum checks for tables that have had a large number of inserted, updated or deleted tuples. These checks use the statistics collection facility; therefore, autovacuum cannot be used unless track_counts is set to true. In the default configuration, autovacuuming is enabled and the related configuration parameters are appropriately set.

The “autovacuum daemon” actually consists of multiple processes. There is a persistent daemon process, called the autovacuum launcher, which is in charge of starting autovacuum worker processes for all databases. The launcher will distribute the work across time, attempting to start one worker within each database every autovacuum_naptime seconds. (Therefore, if the installation has N databases, a new worker will be launched every autovacuum_naptime/N seconds.) A maximum of autovacuum_max_workers worker processes are allowed to run at the same time. If there are more than autovacuum_max_workers databases to be processed, the next database will be processed as soon as the first worker finishes. Each worker process will check each table within its database and execute VACUUM and/or ANALYZE as needed. log_autovacuum_min_duration can be set to monitor autovacuum workers’ activity.

If several large tables all become eligible for vacuuming in a short amount of time, all autovacuum workers might become occupied with vacuuming those tables for a long period. This would result in
other tables and databases not being vacuumed until a worker becomes available. There is no limit on how many workers might be in a single database, but workers do try to avoid repeating work that has already been done by other workers. Note that the number of running workers does not count towards max_connections or superuser_reserved_connections limits.

Tables whose relfrozenxid value is more than autovacuum_freeze_max_age transactions old are always vacuumed (this also applies to those tables whose freeze max age has been modified via storage parameters; see below). Otherwise, if the number of tuples obsoleted since the last VACUUM exceeds the “vacuum threshold”, the table is vacuumed. The vacuum threshold is defined as:

\[
\text{vacuum threshold} = \text{vacuum base threshold} + \text{vacuum scale factor} \times \text{number of tuples}
\]

where the vacuum base threshold is autovacuum_vacuum_threshold, the vacuum scale factor is autovacuum_vacuum_scale_factor, and the number of tuples is pg_class.reltuples. The number of obsolete tuples is obtained from the statistics collector; it is a semi-accurate count updated by each UPDATE and DELETE operation. (It is only semi-accurate because some information might be lost under heavy load.) If the relfrozenxid value of the table is more than vacuum_freeze_table_age transactions old, an aggressive vacuum is performed to freeze old tuples and advance relfrozenxid; otherwise, only pages that have been modified since the last vacuum are scanned.

For analyze, a similar condition is used: the threshold, defined as:

\[
\text{analyze threshold} = \text{analyze base threshold} + \text{analyze scale factor} \times \text{number of tuples}
\]

is compared to the total number of tuples inserted, updated, or deleted since the last ANALYZE.

Temporary tables cannot be accessed by autovacuum. Therefore, appropriate vacuum and analyze operations should be performed via session SQL commands.

The default thresholds and scale factors are taken from postgresql.conf, but it is possible to override them (and many other autovacuum control parameters) on a per-table basis; see Storage Parameters for more information. If a setting has been changed via a table's storage parameters, that value is used when processing that table; otherwise the global settings are used. See Section 19.10 for more details on the global settings.

When multiple workers are running, the autovacuum cost delay parameters (see Section 19.4.4) are “balanced” among all the running workers, so that the total I/O impact on the system is the same regardless of the number of workers actually running. However, any workers processing tables whose per-table autovacuum_vacuum_cost_delay or autovacuum_vacuum_cost_limit storage parameters have been set are not considered in the balancing algorithm.

### 24.2. Routine Reindexing

In some situations it is worthwhile to rebuild indexes periodically with the REINDEX command or a series of individual rebuilding steps.

B-tree index pages that have become completely empty are reclaimed for re-use. However, there is still a possibility of inefficient use of space: if all but a few index keys on a page have been deleted, the page remains allocated. Therefore, a usage pattern in which most, but not all, keys in each range are eventually deleted will see poor use of space. For such usage patterns, periodic reindexing is recommended.

The potential for bloat in non-B-tree indexes has not been well researched. It is a good idea to periodically monitor the index's physical size when using any non-B-tree index type.
Also, for B-tree indexes, a freshly-constructed index is slightly faster to access than one that has been updated many times because logically adjacent pages are usually also physically adjacent in a newly built index. (This consideration does not apply to non-B-tree indexes.) It might be worthwhile to reindex periodically just to improve access speed.

REINDEX can be used safely and easily in all cases. But since the command requires an exclusive table lock, it is often preferable to execute an index rebuild with a sequence of creation and replacement steps. Index types that support CREATE INDEX with the CONCURRENTLY option can instead be recreated that way. If that is successful and the resulting index is valid, the original index can then be replaced by the newly built one using a combination of ALTER INDEX and DROP INDEX. When an index is used to enforce uniqueness or other constraints, ALTER TABLE might be necessary to swap the existing constraint with one enforced by the new index. Review this alternate multistep rebuild approach carefully before using it as there are limitations on which indexes can be reindexed this way, and errors must be handled.

### 24.3. Log File Maintenance

It is a good idea to save the database server's log output somewhere, rather than just discarding it via /dev/null. The log output is invaluable when diagnosing problems. However, the log output tends to be voluminous (especially at higher debug levels) so you won't want to save it indefinitely. You need to rotate the log files so that new log files are started and old ones removed after a reasonable period of time.

If you simply direct the stderr of postgres into a file, you will have log output, but the only way to truncate the log file is to stop and restart the server. This might be acceptable if you are using PostgreSQL in a development environment, but few production servers would find this behavior acceptable.

A better approach is to send the server's stderr output to some type of log rotation program. There is a built-in log rotation facility, which you can use by setting the configuration parameter logging_collector to true in postgresql.conf. The control parameters for this program are described in Section 19.8.1. You can also use this approach to capture the log data in machine readable CSV (comma-separated values) format.

Alternatively, you might prefer to use an external log rotation program if you have one that you are already using with other server software. For example, the rotatelogs tool included in the Apache distribution can be used with PostgreSQL. To do this, just pipe the server's stderr output to the desired program. If you start the server with pg_ctl, then stderr is already redirected to stdout, so you just need a pipe command, for example:

```
pg_ctl start | rotatelogs /var/log/pgsql_log 86400
```

Another production-grade approach to managing log output is to send it to syslog and let syslog deal with file rotation. To do this, set the configuration parameter log_destination to syslog (to log to syslog only) in postgresql.conf. Then you can send a SIGHUP signal to the syslog daemon whenever you want to force it to start writing a new log file. If you want to automate log rotation, the logrotate program can be configured to work with log files from syslog.

On many systems, however, syslog is not very reliable, particularly with large log messages; it might truncate or drop messages just when you need them the most. Also, on Linux, syslog will flush each message to disk, yielding poor performance. (You can use a “-” at the start of the file name in the syslog configuration file to disable syncing.)

Note that all the solutions described above take care of starting new log files at configurable intervals, but they do not handle deletion of old, no-longer-useful log files. You will probably want to set up a batch job to periodically delete old log files. Another possibility is to configure the rotation program so that old log files are overwritten cyclically.
pgBadger\textsuperscript{2} is an external project that does sophisticated log file analysis. check\_postgres\textsuperscript{3} provides Nagios alerts when important messages appear in the log files, as well as detection of many other extraordinary conditions.

\textsuperscript{2} \url{http://dalibo.github.io/pgbadger/}
\textsuperscript{3} \url{https://bucardo.org/check_postgres/}
Chapter 25. Backup and Restore

As with everything that contains valuable data, PostgreSQL databases should be backed up regularly. While the procedure is essentially simple, it is important to have a clear understanding of the underlying techniques and assumptions.

There are three fundamentally different approaches to backing up PostgreSQL data:

- **SQL dump**
- **File system level backup**
- **Continuous archiving**

Each has its own strengths and weaknesses; each is discussed in turn in the following sections.

### 25.1. SQL Dump

The idea behind this dump method is to generate a file with SQL commands that, when fed back to the server, will recreate the database in the same state as it was at the time of the dump. PostgreSQL provides the utility program `pg_dump` for this purpose. The basic usage of this command is:

```
pg_dump dbname > dumpfile
```

As you see, `pg_dump` writes its result to the standard output. We will see below how this can be useful. While the above command creates a text file, `pg_dump` can create files in other formats that allow for parallelism and more fine-grained control of object restoration.

`pg_dump` is a regular PostgreSQL client application (albeit a particularly clever one). This means that you can perform this backup procedure from any remote host that has access to the database. But remember that `pg_dump` does not operate with special permissions. In particular, it must have read access to all tables that you want to back up, so in order to back up the entire database you almost always have to run it as a database superuser. (If you do not have sufficient privileges to back up the entire database, you can still back up portions of the database to which you do have access using options such as `-n schema` or `-t table`.)

To specify which database server `pg_dump` should contact, use the command line options `-h host` and `-p port`. The default host is the local host or whatever your `PGHOST` environment variable specifies. Similarly, the default port is indicated by the `PGPORT` environment variable or, failing that, by the compiled-in default. (Conveniently, the server will normally have the same compiled-in default.)

Like any other PostgreSQL client application, `pg_dump` will by default connect with the database user name that is equal to the current operating system user name. To override this, either specify the `-U` option or set the environment variable `PGUSER`. Remember that `pg_dump` connections are subject to the normal client authentication mechanisms (which are described in Chapter 20).

An important advantage of `pg_dump` over the other backup methods described later is that `pg_dump`'s output can generally be re-loaded into newer versions of PostgreSQL, whereas file-level backups and continuous archiving are both extremely server-version-specific. `pg_dump` is also the only method that will work when transferring a database to a different machine architecture, such as going from a 32-bit to a 64-bit server.

Dumps created by `pg_dump` are internally consistent, meaning, the dump represents a snapshot of the database at the time `pg_dump` began running. `pg_dump` does not block other operations on the database while it is working. (Exceptions are those operations that need to operate with an exclusive lock, such as most forms of `ALTER TABLE`.)
25.1.1. Restoring the Dump

Text files created by `pg_dump` are intended to be read in by the `psql` program. The general command form to restore a dump is

```
psql dbname < dumpfile
```

where `dumpfile` is the file output by the `pg_dump` command. The database `dbname` will not be created by this command, so you must create it yourself from `template0` before executing `psql (e.g., with `createdb -T template0 dbname`). `psql` supports options similar to `pg_dump` for specifying the database server to connect to and the user name to use. See the `psql` reference page for more information. Non-text file dumps are restored using the `pg_restore` utility.

Before restoring an SQL dump, all the users who own objects or were granted permissions on objects in the dumped database must already exist. If they do not, the restore will fail to recreate the objects with the original ownership and/or permissions. (Sometimes this is what you want, but usually it is not.)

By default, the `psql` script will continue to execute after an SQL error is encountered. You might wish to run `psql` with the `ON_ERROR_STOP` variable set to alter that behavior and have `psql` exit with an exit status of 3 if an SQL error occurs:

```
psql --set ON_ERROR_STOP=on dbname < dumpfile
```

Either way, you will only have a partially restored database. Alternately, you can specify that the whole dump should be restored as a single transaction, so the restore is either fully completed or fully rolled back. This mode can be specified by passing the `--single-transaction` command-line options to `psql`. When using this mode, be aware that even a minor error can rollback a restore that has already run for many hours. However, that might still be preferable to manually cleaning up a complex database after a partially restored dump.

The ability of `pg_dump` and `psql` to write to or read from pipes makes it possible to dump a database directly from one server to another, for example:

```
pg_dump -h host1 dbname | psql -h host2 dbname
```

### Important

The dumps produced by `pg_dump` are relative to `template0`. This means that any languages, procedures, etc. added via `template1` will also be dumped by `pg_dump`. As a result, when restoring, if you are using a customized `template1`, you must create the empty database from `template0`, as in the example above.

After restoring a backup, it is wise to run `ANALYZE` on each database so the query optimizer has useful statistics; see Section 24.1.3 and Section 24.1.6 for more information. For more advice on how to load large amounts of data into PostgreSQL efficiently, refer to Section 14.4.

25.1.2. Using `pg_dumpall`

`pg_dump` dumps only a single database at a time, and it does not dump information about roles or tablespaces (because those are cluster-wide rather than per-database). To support convenient dumping of the entire contents of a database cluster, the `pg_dumpall` program is provided. `pg_dumpall` backs up each database in a given cluster, and also preserves cluster-wide data such as role and tablespace definitions. The basic usage of this command is:
Backup and Restore

```
pg_dumpall > dumpfile
```

The resulting dump can be restored with psql:

```
psql -f dumpfile postgres
```

(Actually, you can specify any existing database name to start from, but if you are loading into an empty cluster then `postgres` should usually be used.) It is always necessary to have database superuser access when restoring a pg_dumpall dump, as that is required to restore the role and tablespace information. If you use tablespaces, make sure that the tablespace paths in the dump are appropriate for the new installation.

`pg_dumpall` works by emitting commands to re-create roles, tablespaces, and empty databases, then invoking `pg_dump` for each database. This means that while each database will be internally consistent, the snapshots of different databases are not synchronized.

Cluster-wide data can be dumped alone using the `pg_dumpall --globals-only` option. This is necessary to fully backup the cluster if running the `pg_dump` command on individual databases.

### 25.1.3. Handling Large Databases

Some operating systems have maximum file size limits that cause problems when creating large `pg_dump` output files. Fortunately, `pg_dump` can write to the standard output, so you can use standard Unix tools to work around this potential problem. There are several possible methods:

**Use compressed dumps.** You can use your favorite compression program, for example `gzip`:

```
pg_dump dbname | gzip > filename.gz
```

Reload with:

```
gunzip -c filename.gz | psql dbname
```

or:

```
cat filename.gz | gunzip | psql dbname
```

**Use `split`.** The `split` command allows you to split the output into smaller files that are acceptable in size to the underlying file system. For example, to make chunks of 1 megabyte:

```
pg_dump dbname | split -b 1m - filename
```

Reload with:

```
cat filename* | psql dbname
```

**Use `pg_dump`'s custom dump format.** If PostgreSQL was built on a system with the zlib compression library installed, the custom dump format will compress data as it writes it to the output file. This will produce dump file sizes similar to using `gzip`, but it has the added advantage that tables can be restored selectively. The following command dumps a database using the custom dump format:

```
pg_dump -Fc dbname > filename
```

A custom-format dump is not a script for psql, but instead must be restored with `pg_restore`, for example:
Backup and Restore

`pg_restore -d dbname filename`

See the `pg_dump` and `pg_restore` reference pages for details.

For very large databases, you might need to combine sleep with one of the other two approaches.

**Use `pg_dump`'s parallel dump feature.** To speed up the dump of a large database, you can use `pg_dump`'s parallel mode. This will dump multiple tables at the same time. You can control the degree of parallelism with the `-j` parameter. Parallel dumps are only supported for the "directory" archive format.

`pg_dump -j num -F d -f out.dir dbname`

You can use `pg_restore -j` to restore a dump in parallel. This will work for any archive of either the "custom" or the "directory" archive mode, whether or not it has been created with `pg_dump -j`.

### 25.2. File System Level Backup

An alternative backup strategy is to directly copy the files that PostgreSQL uses to store the data in the database; Section 18.2 explains where these files are located. You can use whatever method you prefer for doing file system backups; for example:

```
tar -cf backup.tar /usr/local/pgsql/data
```

There are two restrictions, however, which make this method impractical, or at least inferior to the `pg_dump` method:

1. The database server must be shut down in order to get a usable backup. Half-way measures such as disallowing all connections will not work (in part because `tar` and similar tools do not take an atomic snapshot of the state of the file system, but also because of internal buffering within the server). Information about stopping the server can be found in Section 18.5. Needless to say, you also need to shut down the server before restoring the data.

2. If you have dug into the details of the file system layout of the database, you might be tempted to try to back up or restore only certain individual tables or databases from their respective files or directories. This will not work because the information contained in these files is not usable without the commit log files, `pg_xact/*`, which contain the commit status of all transactions. A table file is only usable with this information. Of course it is also impossible to restore only a table and the associated `pg_xact` data because that would render all other tables in the database cluster useless. So file system backups only work for complete backup and restoration of an entire database cluster.

An alternative file-system backup approach is to make a “consistent snapshot” of the data directory, if the file system supports that functionality (and you are willing to trust that it is implemented correctly). The typical procedure is to make a “frozen snapshot” of the volume containing the database, then copy the whole data directory (not just parts, see above) from the snapshot to a backup device, then release the frozen snapshot. This will work even while the database server is running. However, a backup created in this way saves the database files in a state as if the database server was not properly shut down; therefore, when you start the database server on the backed-up data, it will think the previous server instance crashed and will replay the WAL log. This is not a problem; just be aware of it (and be sure to include the WAL files in your backup). You can perform a `CHECKPOINT` before taking the snapshot to reduce recovery time.

If your database is spread across multiple file systems, there might not be any way to obtain exactly-simultaneous frozen snapshots of all the volumes. For example, if your data files and WAL log are on different disks, or if tablespaces are on different file systems, it might not be possible to use snapshot backup because the snapshots must be simultaneous. Read your file system documentation very carefully before trusting the consistent-snapshot technique in such situations.
If simultaneous snapshots are not possible, one option is to shut down the database server long enough to establish all the frozen snapshots. Another option is to perform a continuous archiving base backup (Section 25.3.2) because such backups are immune to file system changes during the backup. This requires enabling continuous archiving just during the backup process; restore is done using continuous archive recovery (Section 25.3.4).

Another option is to use rsync to perform a file system backup. This is done by first running rsync while the database server is running, then shutting down the database server long enough to do an `rsync --checksum`. (`--checksum` is necessary because `rsync` only has file modification-time granularity of one second.) The second rsync will be quicker than the first, because it has relatively little data to transfer, and the end result will be consistent because the server was down. This method allows a file system backup to be performed with minimal downtime.

Note that a file system backup will typically be larger than an SQL dump. (pg_dump does not need to dump the contents of indexes for example, just the commands to recreate them.) However, taking a file system backup might be faster.

# 25.3. Continuous Archiving and Point-in-Time Recovery (PITR)

At all times, PostgreSQL maintains a write ahead log (WAL) in the `pg_wal/` subdirectory of the cluster's data directory. The log records every change made to the database's data files. This log exists primarily for crash-safety purposes: if the system crashes, the database can be restored to consistency by “replaying” the log entries made since the last checkpoint. However, the existence of the log makes it possible to use a third strategy for backing up databases: we can combine a file-system-level backup with backup of the WAL files. If recovery is needed, we restore the file system backup and then replay from the backed-up WAL files to bring the system to a current state. This approach is more complex to administer than either of the previous approaches, but it has some significant benefits:

- We do not need a perfectly consistent file system backup as the starting point. Any internal inconsistency in the backup will be corrected by log replay (this is not significantly different from what happens during crash recovery). So we do not need a file system snapshot capability, just tar or a similar archiving tool.

- Since we can combine an indefinitely long sequence of WAL files for replay, continuous backup can be achieved simply by continuing to archive the WAL files. This is particularly valuable for large databases, where it might not be convenient to take a full backup frequently.

- It is not necessary to replay the WAL entries all the way to the end. We could stop the replay at any point and have a consistent snapshot of the database as it was at that time. Thus, this technique supports point-in-time recovery: it is possible to restore the database to its state at any time since your base backup was taken.

- If we continuously feed the series of WAL files to another machine that has been loaded with the same base backup file, we have a warm standby system: at any point we can bring up the second machine and it will have a nearly-current copy of the database.

## Note

pg_dump and pg_dumpall do not produce file-system-level backups and cannot be used as part of a continuous-archiving solution. Such dumps are logical and do not contain enough information to be used by WAL replay.

As with the plain file-system-backup technique, this method can only support restoration of an entire database cluster, not a subset. Also, it requires a lot of archival storage: the base backup might be
backup and a busy system will generate many megabytes of WAL traffic that have to be archived. Still, it is the preferred backup technique in many situations where high reliability is needed.

To recover successfully using continuous archiving (also called “online backup” by many database vendors), you need a continuous sequence of archived WAL files that extends back at least as far as the start time of your backup. So to get started, you should set up and test your procedure for archiving WAL files before you take your first base backup. Accordingly, we first discuss the mechanics of archiving WAL files.

25.3.1. Setting Up WAL Archiving

In an abstract sense, a running PostgreSQL system produces an indefinitely long sequence of WAL records. The system physically divides this sequence into WAL segment files, which are normally 16MB apiece (although the segment size can be altered during initdb). The segment files are given numeric names that reflect their position in the abstract WAL sequence. When not using WAL archiving, the system normally creates just a few segment files and then “recycles” them by renaming no-longer-needed segment files to higher segment numbers. It's assumed that segment files whose contents precede the last checkpoint are no longer of interest and can be recycled.

When archiving WAL data, we need to capture the contents of each segment file once it is filled, and save that data somewhere before the segment file is recycled for reuse. Depending on the application and the available hardware, there could be many different ways of “saving the data somewhere”: we could copy the segment files to an NFS-mounted directory on another machine, write them onto a tape drive (ensuring that you have a way of identifying the original name of each file), or batch them together and burn them onto CDs, or something else entirely. To provide the database administrator with flexibility, PostgreSQL tries not to make any assumptions about how the archiving will be done. Instead, PostgreSQL lets the administrator specify a shell command to be executed to copy a completed segment file to wherever it needs to go. The command could be as simple as a cp, or it could invoke a complex shell script — it's all up to you.

To enable WAL archiving, set the wal_level configuration parameter to replica or higher, archive_mode to on, and specify the shell command to use in the archive_command configuration parameter. In practice these settings will always be placed in the postgresql.conf file. In archive_command, %p is replaced by the path name of the file to archive, while %f is replaced by only the file name. (The path name is relative to the current working directory, i.e., the cluster's data directory.) Use %% if you need to embed an actual % character in the command. The simplest useful command is something like:

```
archive_command = 'test ! -f /mnt/server/archivedir/%f && cp %p /mnt/server/archivedir/%f'  # Unix
archive_command = 'copy "%p" "C:\server\archivedir\%f"'  # Windows
```

which will copy archivable WAL segments to the directory /mnt/server/archivedir. (This is an example, not a recommendation, and might not work on all platforms.) After the %p and %f parameters have been replaced, the actual command executed might look like this:

```
test ! -f /mnt/server/archivedir/00000001000000A900000065 && cp pg_wal/00000001000000A900000065 /mnt/server/archivedir/00000001000000A900000065
```

A similar command will be generated for each new file to be archived.

The archive command will be executed under the ownership of the same user that the PostgreSQL server is running as. Since the series of WAL files being archived contains effectively everything in your database, you will want to be sure that the archived data is protected from prying eyes; for example, archive into a directory that does not have group or world read access.
It is important that the archive command return zero exit status if and only if it succeeds. Upon getting a zero result, PostgreSQL will assume that the file has been successfully archived, and will remove or recycle it. However, a nonzero status tells PostgreSQL that the file was not archived; it will try again periodically until it succeeds.

The archive command should generally be designed to refuse to overwrite any pre-existing archive file. This is an important safety feature to preserve the integrity of your archive in case of administrator error (such as sending the output of two different servers to the same archive directory).

It is advisable to test your proposed archive command to ensure that it indeed does not overwrite an existing file, and that it returns nonzero status in this case. The example command above for Unix ensures this by including a separate test step. On some Unix platforms, $cp$ has switches such as $-i$ that can be used to do the same thing less verbosely, but you should not rely on these without verifying that the right exit status is returned. (In particular, GNU $cp$ will return status zero when $-i$ is used and the target file already exists, which is not the desired behavior.)

While designing your archiving setup, consider what will happen if the archive command fails repeatedly because some aspect requires operator intervention or the archive runs out of space. For example, this could occur if you write to tape without an autochanger; when the tape fills, nothing further can be archived until the tape is swapped. You should ensure that any error condition or request to a human operator is reported appropriately so that the situation can be resolved reasonably quickly. The $pg_wal/$ directory will continue to fill with WAL segment files until the situation is resolved. (If the file system containing $pg_wal/$ fills up, PostgreSQL will do a PANIC shutdown. No committed transactions will be lost, but the database will remain offline until you free some space.)

The speed of the archiving command is unimportant as long as it can keep up with the average rate at which your server generates WAL data. Normal operation continues even if the archiving process falls a little behind. If archiving falls significantly behind, this will increase the amount of data that would be lost in the event of a disaster. It will also mean that the $pg_wal/$ directory will contain large numbers of not-yet-archived segment files, which could eventually exceed available disk space. You are advised to monitor the archiving process to ensure that it is working as you intend.

In writing your archive command, you should assume that the file names to be archived can be up to 64 characters long and can contain any combination of ASCII letters, digits, and dots. It is not necessary to preserve the original relative path ($%p$) but it is necessary to preserve the file name ($%f$).

Note that although WAL archiving will allow you to restore any modifications made to the data in your PostgreSQL database, it will not restore changes made to configuration files (that is, $postgresql.conf$, $pg_hba.conf$ and $pg_ident.conf$), since those are edited manually rather than through SQL operations. You might wish to keep the configuration files in a location that will be backed up by your regular file system backup procedures. See Section 19.2 for how to relocate the configuration files.

The archive command is only invoked on completed WAL segments. Hence, if your server generates only little WAL traffic (or has slack periods where it does so), there could be a long delay between the completion of a transaction and its safe recording in archive storage. To put a limit on how old unarchived data can be, you can set archive_timeout to force the server to switch to a new WAL segment file at least that often. Note that archived files that are archived early due to a forced switch are still the same length as completely full files. It is therefore unwise to set a very short archive_timeout — it will bloat your archive storage. archive_timeout settings of a minute or so are usually reasonable.

Also, you can force a segment switch manually with $pg_switch_wal$ if you want to ensure that a just-finished transaction is archived as soon as possible. Other utility functions related to WAL management are listed in Table 9.79.

When wal_level is minimal some SQL commands are optimized to avoid WAL logging, as described in Section 14.4.7. If archiving or streaming replication were turned on during execution of one of these statements, WAL would not contain enough information for archive recovery. (Crash recovery is unaffected.) For this reason, wal_level can only be changed at server start. However,
archive_command can be changed with a configuration file reload. If you wish to temporarily stop
archiving, one way to do it is to set archive_command to the empty string ("""). This will cause
WAL files to accumulate in pg_wal/ until a working archive_command is re-established.

25.3.2. Making a Base Backup

The easiest way to perform a base backup is to use the pg_basebackup tool. It can create a base backup
either as regular files or as a tar archive. If more flexibility than pg_basebackup can provide is required,
you can also make a base backup using the low level API (see Section 25.3.3).

It is not necessary to be concerned about the amount of time it takes to make a base backup. How-
ever, if you normally run the server with full_page_writes disabled, you might notice a drop
in performance while the backup runs since full_page_writes is effectively forced on during
backup mode.

To make use of the backup, you will need to keep all the WAL segment files generated during
and after the file system backup. To aid you in doing this, the base backup process creates a backup
history file that is immediately stored into the WAL archive area. This file is named after the
first WAL segment file that you need for the file system backup. For example, if the starting WAL
file is 0000000100001234000055CD the backup history file will be named something like
0000000100001234000055CD.007C9330.backup. (The second part of the file name stands
for an exact position within the WAL file, and can ordinarily be ignored.) Once you have safely
archived the file system backup and the WAL segment files used during the backup (as specified in the
backup history file), all archived WAL segments with names numerically less are no longer needed
to recover the file system backup and can be deleted. However, you should consider keeping several
backup sets to be absolutely certain that you can recover your data.

The backup history file is just a small text file. It contains the label string you gave to pg_basebackup,
as well as the starting and ending times and WAL segments of the backup. If you used the label to
identify the associated dump file, then the archived history file is enough to tell you which dump file
to restore.

Since you have to keep around all the archived WAL files back to your last base backup, the interval
between base backups should usually be chosen based on how much storage you want to expend on
archived WAL files. You should also consider how long you are prepared to spend recovering, if
recovery should be necessary — the system will have to replay all those WAL segments, and that
could take awhile if it has been a long time since the last base backup.

25.3.3. Making a Base Backup Using the Low Level API

The procedure for making a base backup using the low level APIs contains a few more steps than the
pg_basebackup method, but is relatively simple. It is very important that these steps are executed in
sequence, and that the success of a step is verified before proceeding to the next step.

Low level base backups can be made in a non-exclusive or an exclusive way. The non-exclusive
method is recommended and the exclusive one is deprecated and will eventually be removed.

25.3.3.1. Making a non-exclusive low level backup

A non-exclusive low level backup is one that allows other concurrent backups to be running (both
those started using the same backup API and those started using pg_basebackup).

1. Ensure that WAL archiving is enabled and working.

2. Connect to the server (it does not matter which database) as a user with rights to run pg_start_backup
   (superuser, or a user who has been granted EXECUTE on the function) and issue the command:

   SELECT pg_start_backup('label', false, false);
where \texttt{label} is any string you want to use to uniquely identify this backup operation. The connection calling \texttt{pg\_start\_backup} must be maintained until the end of the backup, or the backup will be automatically aborted.

By default, \texttt{pg\_start\_backup} can take a long time to finish. This is because it performs a checkpoint, and the I/O required for the checkpoint will be spread out over a significant period of time, by default half your inter-checkpoint interval (see the configuration parameter \texttt{checkpoint\_completion\_target}). This is usually what you want, because it minimizes the impact on query processing. If you want to start the backup as soon as possible, change the second parameter to \texttt{true}, which will issue an immediate checkpoint using as much I/O as available.

The third parameter being \texttt{false} tells \texttt{pg\_start\_backup} to initiate a non-exclusive base backup.

3. Perform the backup, using any convenient file-system-backup tool such as tar or cpio (not \texttt{pg\_dump} or \texttt{pg\_dumpall}). It is neither necessary nor desirable to stop normal operation of the database while you do this. See Section 25.3.3.3 for things to consider during this backup.

4. In the same connection as before, issue the command:

\begin{verbatim}
SELECT * FROM pg_stop_backup(false, true);
\end{verbatim}

This terminates backup mode. On a primary, it also performs an automatic switch to the next WAL segment. On a standby, it is not possible to automatically switch WAL segments, so you may wish to run \texttt{pg\_switch\_wal} on the primary to perform a manual switch. The reason for the switch is to arrange for the last WAL segment file written during the backup interval to be ready to archive.

The \texttt{pg\_stop\_backup} will return one row with three values. The second of these fields should be written to a file named \texttt{backup\_label} in the root directory of the backup. The third field should be written to a file named \texttt{tablespace\_map} unless the field is empty. These files are vital to the backup working, and must be written without modification.

5. Once the WAL segment files active during the backup are archived, you are done. The file identified by \texttt{pg\_stop\_backup}’s first return value is the last segment that is required to form a complete set of backup files. On a primary, if \texttt{archive\_mode} is enabled and the \texttt{wait\_for\_archive} parameter is \texttt{true}, \texttt{pg\_stop\_backup} does not return until the last segment has been archived. On a standby, \texttt{archive\_mode} must be \texttt{always} in order for \texttt{pg\_stop\_backup} to wait. Archiving of these files happens automatically since you have already configured \texttt{archive\_command}. In most cases this happens quickly, but you are advised to monitor your archive system to ensure there are no delays. If the archive process has fallen behind because of failures of the archive command, it will keep retrying until the archive succeeds and the backup is complete. If you wish to place a time limit on the execution of \texttt{pg\_stop\_backup}, set an appropriate \texttt{statement\_timeout} value, but make note that if \texttt{pg\_stop\_backup} terminates because of this your backup may not be valid.

If the backup process monitors and ensures that all WAL segment files required for the backup are successfully archived then the \texttt{wait\_for\_archive} parameter (which defaults to \texttt{true}) can be set to \texttt{false} to have \texttt{pg\_stop\_backup} return as soon as the stop backup record is written to the WAL. By default, \texttt{pg\_stop\_backup} will wait until all WAL has been archived, which can take some time. This option must be used with caution: if WAL archiving is not monitored correctly then the backup might not include all of the WAL files and will therefore be incomplete and not able to be restored.

25.3.3.2. Making an exclusive low level backup

The process for an exclusive backup is mostly the same as for a non-exclusive one, but it differs in a few key steps. This type of backup can only be taken on a primary and does not allow concurrent backups. Prior to PostgreSQL 9.6, this was the only low-level method available, but it is now recommended that all users upgrade their scripts to use non-exclusive backups if possible.
1. Ensure that WAL archiving is enabled and working.

2. Connect to the server (it does not matter which database) as a user with rights to run `pg_start_backup` (superuser, or a user who has been granted EXECUTE on the function) and issue the command:

   ```sql
   SELECT pg_start_backup('label');
   ``

   where `label` is any string you want to use to uniquely identify this backup operation. `pg_start_backup` creates a backup label file, called `backup_label`, in the cluster directory with information about your backup, including the start time and label string. The function also creates a tablespace map file, called `tablespace_map`, in the cluster directory with information about tablespace symbolic links in `pg_tblspc/` if one or more such link is present. Both files are critical to the integrity of the backup, should you need to restore from it.

   By default, `pg_start_backup` can take a long time to finish. This is because it performs a checkpoint, and the I/O required for the checkpoint will be spread out over a significant period of time, by default half your inter-checkpoint interval (see the configuration parameter `checkpoint_completion_target`). This is usually what you want, because it minimizes the impact on query processing. If you want to start the backup as soon as possible, use:

   ```sql
   SELECT pg_start_backup('label', true);
   ``

   This forces the checkpoint to be done as quickly as possible.

3. Perform the backup, using any convenient file-system-backup tool such as tar or cpio (not `pg_dump` or `pg_dumpall`). It is neither necessary nor desirable to stop normal operation of the database while you do this. See Section 25.3.3.3 for things to consider during this backup.

   Note that if the server crashes during the backup it may not be possible to restart until the backup_label file has been manually deleted from the `PGDATA` directory.

4. Again connect to the database as a user with rights to run `pg_stop_backup` (superuser, or a user who has been granted EXECUTE on the function), and issue the command:

   ```sql
   SELECT pg_stop_backup();
   ``

   This function terminates backup mode and performs an automatic switch to the next WAL segment.

   The reason for the switch is to arrange for the last WAL segment written during the backup interval to be ready to archive.

5. Once the WAL segment files active during the backup are archived, you are done. The file identified by `pg_stop_backup`'s result is the last segment that is required to form a complete set of backup files. If `archive_mode` is enabled, `pg_stop_backup` does not return until the last segment has been archived. Archiving of these files happens automatically since you have already configured `archive_command`. In most cases this happens quickly, but you are advised to monitor your archive system to ensure there are no delays. If the archive process has fallen behind because of failures of the archive command, it will keep retrying until the archive succeeds and the backup is complete. If you wish to place a time limit on the execution of `pg_stop_backup`, set an appropriate `statement_timeout` value, but make note that if `pg_stop_backup` terminates because of this your backup may not be valid.

25.3.3.3. Backing up the data directory

Some file system backup tools emit warnings or errors if the files they are trying to copy change while the copy proceeds. When taking a base backup of an active database, this situation is normal and not an error. However, you need to ensure that you can distinguish complaints of this sort from real errors. For example, some versions of rsync return a separate exit code for “vanished source files”, and you
Backup and Restore

can write a driver script to accept this exit code as a non-error case. Also, some versions of GNU tar return an error code indistinguishable from a fatal error if a file was truncated while tar was copying it. Fortunately, GNU tar versions 1.16 and later exit with 1 if a file was changed during the backup, and 2 for other errors. With GNU tar version 1.23 and later, you can use the warning options --warning=no-file-changed --warning=no-file-removed to hide the related warning messages.

Be certain that your backup includes all of the files under the database cluster directory (e.g., /usr/local/pgsql/data). If you are using tablespaces that do not reside underneath this directory, be careful to include them as well (and be sure that your backup archives symbolic links as links, otherwise the restore will corrupt your tablespaces).

You should, however, omit from the backup the files within the cluster's pg_wal/ subdirectory. This slight adjustment is worthwhile because it reduces the risk of mistakes when restoring. This is easy to arrange if pg_wal/ is a symbolic link pointing to someplace outside the cluster directory, which is a common setup anyway for performance reasons. You might also want to exclude postmaster.pid and postmaster.opts, which record information about the running postmaster, not about the postmaster which will eventually use this backup. (These files can confuse pg_ctl.)

It is often a good idea to also omit from the backup the files within the cluster's pg_replslot/ directory, so that replication slots that exist on the master do not become part of the backup. Otherwise, the subsequent use of the backup to create a standby may result in indefinite retention of WAL files on the standby, and possibly bloat on the master if hot standby feedback is enabled, because the clients that are using those replication slots will still be connecting to and updating the slots on the master, not the standby. Even if the backup is only intended for use in creating a new master, copying the replication slots isn't expected to be particularly useful, since the contents of those slots will likely be badly out of date by the time the new master comes on line.

The contents of the directories pg_dynshmem/, pg_notify/, pg_serial/, pg_snapshots/, pg_stat_tmp/, and pg_subtrans/ (but not the directories themselves) can be omitted from the backup as they will be initialized on postmaster startup. If stats_temp_directory is set and is under the data directory then the contents of that directory can also be omitted.

Any file or directory beginning with pgsql_tmp can be omitted from the backup. These files are removed on postmaster start and the directories will be recreated as needed.

pg_internal.init files can be omitted from the backup whenever a file of that name is found. These files contain relation cache data that is always rebuilt when recovering.

The backup label file includes the label string you gave to pg_start_backup, as well as the time at which pg_start_backup was run, and the name of the starting WAL file. In case of confusion it is therefore possible to look inside a backup file and determine exactly which backup session the dump file came from. The tablespace map file includes the symbolic link names as they exist in the directory pg_tblspc/ and the full path of each symbolic link. These files are not merely for your information; their presence and contents are critical to the proper operation of the system’s recovery process.

It is also possible to make a backup while the server is stopped. In this case, you obviously cannot use pg_start_backup or pg_stop_backup, and you will therefore be left to your own devices to keep track of which backup is which and how far back the associated WAL files go. It is generally better to follow the continuous archiving procedure above.

25.3.4. Recovering Using a Continuous Archive Backup

Okay, the worst has happened and you need to recover from your backup. Here is the procedure:

1. Stop the server, if it’s running.

2. If you have the space to do so, copy the whole cluster data directory and any tablespaces to a temporary location in case you need them later. Note that this precaution will require that you have
enough free space on your system to hold two copies of your existing database. If you do not have
enough space, you should at least save the contents of the cluster's pg_wal subdirectory, as it
might contain logs which were not archived before the system went down.

3. Remove all existing files and subdirectories under the cluster data directory and under the root
directories of any tablespaces you are using.

4. Restore the database files from your file system backup. Be sure that they are restored with the right
ownership (the database system user, not root!) and with the right permissions. If you are using
tablespaces, you should verify that the symbolic links in pg_tblspc/ were correctly restored.

5. Remove any files present in pg_wal/; these came from the file system backup and are therefore
probably obsolete rather than current. If you didn't archive pg_wal/ at all, then recreate it with
proper permissions, being careful to ensure that you re-establish it as a symbolic link if you had
it set up that way before.

6. If you have unarchived WAL segment files that you saved in step 2, copy them into pg_wal/.
   (It is best to copy them, not move them, so you still have the unmodified files if a problem occurs
   and you have to start over.)

7. Create a recovery command file recovery.conf in the cluster data directory (see Chapter 27).
   You might also want to temporarily modify pg_hba.conf to prevent ordinary users from con-
   necting until you are sure the recovery was successful.

8. Start the server. The server will go into recovery mode and proceed to read through the archived
   WAL files it needs. Should the recovery be terminated because of an external error, the server can
simply be restarted and it will continue recovery. Upon completion of the recovery process, the
server will rename recovery.conf to recovery.done (to prevent accidentally re-entering
recovery mode later) and then commence normal database operations.

9. Inspect the contents of the database to ensure you have recovered to the desired state. If not, return
to step 1. If all is well, allow your users to connect by restoring pg_hba.conf to normal.

The key part of all this is to set up a recovery configuration file that describes how you want to recover
and how far the recovery should run. You can use recovery.conf.sample (normally located in
the installation's share/ directory) as a prototype. The one thing that you absolutely must specify
in recovery.conf is the restore_command, which tells PostgreSQL how to retrieve archived
WAL file segments. Like the archive_command, this is a shell command string. It can contain
%f, which is replaced by the name of the desired log file, and %p, which is replaced by the path name
to copy the log file to. (The path name is relative to the current working directory, i.e., the cluster's
data directory.) Write % if you need to embed an actual % character in the command. The simplest
useful command is something like:

```
restore_command = 'cp /mnt/server/archivedir/%f %p'
```

which will copy previously archived WAL segments from the directory /mnt/server/archivedir. Of course, you can use something much more complicated, perhaps even a shell
script that requests the operator to mount an appropriate tape.

It is important that the command return nonzero exit status on failure. The command will be called
requesting files that are not present in the archive; it must return nonzero when so asked. This is not an
error condition. An exception is that if the command was terminated by a signal (other than SIGTERM,
which is used as part of a database server shutdown) or an error by the shell (such as command not
found), then recovery will abort and the server will not start up.

Not all of the requested files will be WAL segment files; you should also expect requests for files with
a suffix of .history. Also be aware that the base name of the %p path will be different from %f;
do not expect them to be interchangeable.
WAL segments that cannot be found in the archive will be sought in `pg_wal/`; this allows use of recent un-archived segments. However, segments that are available from the archive will be used in preference to files in `pg_wal/`.

Normally, recovery will proceed through all available WAL segments, thereby restoring the database to the current point in time (or as close as possible given the available WAL segments). Therefore, a normal recovery will end with a “file not found” message, the exact text of the error message depending upon your choice of `restore_command`. You may also see an error message at the start of recovery for a file named something like `00000001.history`. This is also normal and does not indicate a problem in simple recovery situations; see Section 25.3.5 for discussion.

If you want to recover to some previous point in time (say, right before the junior DBA dropped your main transaction table), just specify the required stopping point in `recovery.conf`. You can specify the stop point, known as the “recovery target”, either by date/time, named restore point or by completion of a specific transaction ID. As of this writing only the date/time and named restore point options are very usable, since there are no tools to help you identify with any accuracy which transaction ID to use.

### Note

The stop point must be after the ending time of the base backup, i.e., the end time of `pg_stop_backup`. You cannot use a base backup to recover to a time when that backup was in progress. (To recover to such a time, you must go back to your previous base backup and roll forward from there.)

If recovery finds corrupted WAL data, recovery will halt at that point and the server will not start. In such a case the recovery process could be re-run from the beginning, specifying a “recovery target” before the point of corruption so that recovery can complete normally. If recovery fails for an external reason, such as a system crash or if the WAL archive has become inaccessible, then the recovery can simply be restarted and it will restart almost from where it failed. Recovery restart works much like checkpointing in normal operation: the server periodically forces all its state to disk, and then updates the `pg_control` file to indicate that the already-processed WAL data need not be scanned again.

## 25.3.5. Timelines

The ability to restore the database to a previous point in time creates some complexities that are akin to science-fiction stories about time travel and parallel universes. For example, in the original history of the database, suppose you dropped a critical table at 5:15PM on Tuesday evening, but didn't realize your mistake until Wednesday noon. Unfazed, you get out your backup, restore to the point-in-time 5:14PM Tuesday evening, and are up and running. In this history of the database universe, you never dropped the table. But suppose you later realize this wasn't such a great idea, and would like to return to sometime Wednesday morning in the original history. You won't be able to if, while your database was up-and-running, it overwrote some of the WAL segment files that led up to the time you now wish you could get back to. Thus, to avoid this, you need to distinguish the series of WAL records generated after you've done a point-in-time recovery from those that were generated in the original database history.

To deal with this problem, PostgreSQL has a notion of **timelines**. Whenever an archive recovery completes, a new timeline is created to identify the series of WAL records generated after that recovery. The timeline ID number is part of WAL segment file names so a new timeline does not overwrite the WAL data generated by previous timelines. It is in fact possible to archive many different timelines. While that might seem like a useless feature, it's often a lifesaver. Consider the situation where you aren't quite sure what point-in-time to recover to, and so have to do several point-in-time recoveries by trial and error until you find the best place to branch off from the old history. Without timelines this process would soon generate an unmanageable mess. With timelines, you can recover to any prior state, including states in timeline branches that you abandoned earlier.
Every time a new timeline is created, PostgreSQL creates a “timeline history” file that shows which timeline it branched off from and when. These history files are necessary to allow the system to pick the right WAL segment files when recovering from an archive that contains multiple timelines. Therefore, they are archived into the WAL archive area just like WAL segment files. The history files are just small text files, so it's cheap and appropriate to keep them around indefinitely (unlike the segment files which are large). You can, if you like, add comments to a history file to record your own notes about how and why this particular timeline was created. Such comments will be especially valuable when you have a thicket of different timelines as a result of experimentation.

The default behavior of recovery is to recover along the same timeline that was current when the base backup was taken. If you wish to recover into some child timeline (that is, you want to return to some state that was itself generated after a recovery attempt), you need to specify the target timeline ID in recovery.conf. You cannot recover into timelines that branched off earlier than the base backup.

25.3.6. Tips and Examples

Some tips for configuring continuous archiving are given here.

25.3.6.1. Standalone Hot Backups

It is possible to use PostgreSQL's backup facilities to produce standalone hot backups. These are backups that cannot be used for point-in-time recovery, yet are typically much faster to backup and restore than pg_dump dumps. (They are also much larger than pg_dump dumps, so in some cases the speed advantage might be negated.)

As with base backups, the easiest way to produce a standalone hot backup is to use the pg_basebackup tool. If you include the -X parameter when calling it, all the write-ahead log required to use the backup will be included in the backup automatically, and no special action is required to restore the backup.

If more flexibility in copying the backup files is needed, a lower level process can be used for standalone hot backups as well. To prepare for low level standalone hot backups, make sure wal_level is set to replica or higher, archive_mode to on, and set up an archive_command that performs archiving only when a switch file exists. For example:

```
archive_command = 'test ! -f /var/lib/pgsql/backup_in_progress
|| (test ! -f /var/lib/pgsql/archive/%f && cp %p /var/lib/pgsql/archive/%f)'
```

This command will perform archiving when /var/lib/pgsql/backup_in_progress exists, and otherwise silently return zero exit status (allowing PostgreSQL to recycle the unwanted WAL file).

With this preparation, a backup can be taken using a script like the following:

```
touch /var/lib/pgsql/backup_in_progress
psql -c "select pg_start_backup('hot_backup');"
tar -cf /var/lib/pgsql/backup.tar /var/lib/pgsql/data/
psql -c "select pg_stop_backup();"
rm /var/lib/pgsql/backup_in_progress
tar -rf /var/lib/pgsql/backup.tar /var/lib/pgsql/archive/
```

The switch file /var/lib/pgsql/backup_in_progress is created first, enabling archiving of completed WAL files to occur. After the backup the switch file is removed. Archived WAL files are then added to the backup so that both base backup and all required WAL files are part of the same tar file. Please remember to add error handling to your backup scripts.

25.3.6.2. Compressed Archive Logs

If archive storage size is a concern, you can use gzip to compress the archive files:
**Backup and Restore**

archive_command = 'gzip < %p > /var/lib/pgsql/archive/%f'

You will then need to use gunzip during recovery:

restore_command = 'gunzip < /mnt/server/archivedir/%f > %p'

### 25.3.6.3. **archive_command** Scripts

Many people choose to use scripts to define their archive command, so that their postgresql.conf entry looks very simple:

archive_command = 'local_backup_script.sh "%p" "%f"'

Using a separate script file is advisable any time you want to use more than a single command in the archiving process. This allows all complexity to be managed within the script, which can be written in a popular scripting language such as bash or perl.

Examples of requirements that might be solved within a script include:

- Copying data to secure off-site data storage
- Batching WAL files so that they are transferred every three hours, rather than one at a time
- Interfacing with other backup and recovery software
- Interfacing with monitoring software to report errors

**Tip**

When using an **archive_command** script, it's desirable to enable logging_collector. Any messages written to stderr from the script will then appear in the database server log, allowing complex configurations to be diagnosed easily if they fail.

### 25.3.7. **Caveats**

At this writing, there are several limitations of the continuous archiving technique. These will probably be fixed in future releases:

- If a CREATE DATABASE command is executed while a base backup is being taken, and then the template database that the CREATE DATABASE copied is modified while the base backup is still in progress, it is possible that recovery will cause those modifications to be propagated into the created database as well. This is of course undesirable. To avoid this risk, it is best not to modify any template databases while taking a base backup.

- CREATE TABLESPACE commands are WAL-logged with the literal absolute path, and will therefore be replayed as tablespace creations with the same absolute path. This might be undesirable if the log is being replayed on a different machine. It can be dangerous even if the log is being replayed on the same machine, but into a new data directory: the replay will still overwrite the contents of the original tablespace. To avoid potential gotchas of this sort, the best practice is to take a new base backup after creating or dropping tablespaces.

It should also be noted that the default WAL format is fairly bulky since it includes many disk page snapshots. These page snapshots are designed to support crash recovery, since we might need to fix partially-written disk pages. Depending on your system hardware and software, the risk of partial writes might be small enough to ignore, in which case you can significantly reduce the total volume of
archived logs by turning off page snapshots using the `full_page_writes` parameter. (Read the notes and
warnings in Chapter 30 before you do so.) Turning off page snapshots does not prevent use of the logs
for PITR operations. An area for future development is to compress archived WAL data by removing
unnecessary page copies even when `full_page_writes` is on. In the meantime, administrators
might wish to reduce the number of page snapshots included in WAL by increasing the checkpoint
interval parameters as much as feasible.
Database servers can work together to allow a second server to take over quickly if the primary server fails (high availability), or to allow several computers to serve the same data (load balancing). Ideally, database servers could work together seamlessly. Web servers serving static web pages can be combined quite easily by merely load-balancing web requests to multiple machines. In fact, read-only database servers can be combined relatively easily too. Unfortunately, most database servers have a read/write mix of requests, and read/write servers are much harder to combine. This is because though read-only data needs to be placed on each server only once, a write to any server has to be propagated to all servers so that future read requests to those servers return consistent results.

This synchronization problem is the fundamental difficulty for servers working together. Because there is no single solution that eliminates the impact of the sync problem for all use cases, there are multiple solutions. Each solution addresses this problem in a different way, and minimizes its impact for a specific workload.

Some solutions deal with synchronization by allowing only one server to modify the data. Servers that can modify data are called read/write, master or primary servers. Servers that track changes in the master are called standby or secondary servers. A standby server that cannot be connected to until it is promoted to a master server is called a warm standby server, and one that can accept connections and serves read-only queries is called a hot standby server.

Some solutions are synchronous, meaning that a data-modifying transaction is not considered committed until all servers have committed the transaction. This guarantees that a failover will not lose any data and that all load-balanced servers will return consistent results no matter which server is queried. In contrast, asynchronous solutions allow some delay between the time of a commit and its propagation to the other servers, opening the possibility that some transactions might be lost in the switch to a backup server, and that load balanced servers might return slightly stale results. Asynchronous communication is used when synchronous would be too slow.

Solutions can also be categorized by their granularity. Some solutions can deal only with an entire database server, while others allow control at the per-table or per-database level.

Performance must be considered in any choice. There is usually a trade-off between functionality and performance. For example, a fully synchronous solution over a slow network might cut performance by more than half, while an asynchronous one might have a minimal performance impact.

The remainder of this section outlines various failover, replication, and load balancing solutions.

### 26.1. Comparison of Different Solutions

#### Shared Disk Failover

Shared disk failover avoids synchronization overhead by having only one copy of the database. It uses a single disk array that is shared by multiple servers. If the main database server fails, the standby server is able to mount and start the database as though it were recovering from a database crash. This allows rapid failover with no data loss.

Shared hardware functionality is common in network storage devices. Using a network file system is also possible, though care must be taken that the file system has full POSIX behavior (see Section 18.2.2). One significant limitation of this method is that if the shared disk array fails or becomes corrupt, the primary and standby servers are both nonfunctional. Another issue is that the standby server should never access the shared storage while the primary server is running.
File System (Block Device) Replication

A modified version of shared hardware functionality is file system replication, where all changes to a file system are mirrored to a file system residing on another computer. The only restriction is that the mirroring must be done in a way that ensures the standby server has a consistent copy of the file system — specifically, writes to the standby must be done in the same order as those on the master. DRBD is a popular file system replication solution for Linux.

Write-Ahead Log Shipping

Warm and hot standby servers can be kept current by reading a stream of write-ahead log (WAL) records. If the main server fails, the standby contains almost all of the data of the main server, and can be quickly made the new master database server. This can be synchronous or asynchronous and can only be done for the entire database server.

A standby server can be implemented using file-based log shipping (Section 26.2) or streaming replication (see Section 26.2.5), or a combination of both. For information on hot standby, see Section 26.5.

Logical Replication

Logical replication allows a database server to send a stream of data modifications to another server. PostgreSQL logical replication constructs a stream of logical data modifications from the WAL. Logical replication allows the data changes from individual tables to be replicated. Logical replication doesn't require a particular server to be designated as a master or a replica but allows data to flow in multiple directions. For more information on logical replication, see Chapter 31. Through the logical decoding interface (Chapter 49), third-party extensions can also provide similar functionality.

Trigger-Based Master-Standby Replication

A master-stable replication setup sends all data modification queries to the master server. The master server asynchronously sends data changes to the standby server. The standby can answer read-only queries while the master server is running. The standby server is ideal for data warehouse queries.

Slony-I is an example of this type of replication, with per-table granularity, and support for multiple standby servers. Because it updates the standby server asynchronously (in batches), there is possible data loss during fail over.

Statement-Based Replication Middleware

With statement-based replication middleware, a program intercepts every SQL query and sends it to one or all servers. Each server operates independently. Read-write queries must be sent to all servers, so that every server receives any changes. But read-only queries can be sent to just one server, allowing the read workload to be distributed among them.

If queries are simply broadcast unmodified, functions like \texttt{random()}, \texttt{CURRENT\_TIMESTAMP}, and sequences can have different values on different servers. This is because each server operates independently, and because SQL queries are broadcast (and not actual modified rows). If this is unacceptable, either the middleware or the application must query such values from a single server and then use those values in write queries. Another option is to use this replication option with a traditional master-stable setup, i.e. data modification queries are sent only to the master and are propagated to the standby servers via master-stable replication, not by the replication middleware. Care must also be taken that all transactions either commit or abort on all servers, perhaps using two-phase commit (\texttt{PREPARE TRANSACTION} and \texttt{COMMIT PREPARED}). Pgpool-II and Continuent Tungsten are examples of this type of replication.

Asynchronous Multimaster Replication

For servers that are not regularly connected, like laptops or remote servers, keeping data consistent among servers is a challenge. Using asynchronous multimaster replication, each server works
High Availability, Load Balancing, and Replication

independently, and periodically communicates with the other servers to identify conflicting transactions. The conflicts can be resolved by users or conflict resolution rules. Bucardo is an example of this type of replication.

Synchronous Multimaster Replication

In synchronous multimaster replication, each server can accept write requests, and modified data is transmitted from the original server to every other server before each transaction commits. Heavy write activity can cause excessive locking, leading to poor performance. In fact, write performance is often worse than that of a single server. Read requests can be sent to any server. Some implementations use shared disk to reduce the communication overhead. Synchronous multimaster replication is best for mostly read workloads, though its big advantage is that any server can accept write requests — there is no need to partition workloads between master and standby servers, and because the data changes are sent from one server to another, there is no problem with non-deterministic functions like \texttt{random()}. PostgreSQL does not offer this type of replication, though PostgreSQL two-phase commit (\texttt{PREPARE TRANSACTION} and \texttt{COMMIT PREPARED}) can be used to implement this in application code or middleware.

Commercial Solutions

Because PostgreSQL is open source and easily extended, a number of companies have taken PostgreSQL and created commercial closed-source solutions with unique failover, replication, and load balancing capabilities.

Table 26.1 summarizes the capabilities of the various solutions listed above.

| Table 26.1. High Availability, Load Balancing, and Replication Feature Matrix |
|--------------------------------------------------|--|--|--|--|--|--|--|--|
| Feature                                         | Shared Disk Failover | File System Replication | Write-Ahead Log Shipping | Logical Replication | Trigger-Based Master-Standby Replication | Statement-Based Replication Middleware | Asynchronous Multimaster Replication | Synchronous Multimaster Replication |
| Most common implementations                     | NAS | DRBD | built-in streaming replication | built-in logical replication, pglogical | Londiste, Slony | pgpool-II | Bucardo |
| Communication method                            | shared disk | disk blocks | WAL | logical decoding | table rows | SQL | table rows | table rows and row locks |
| No special hardware required                    | • | • | • | • | • | • | • |
| Allows multiple master servers                  | • | • | • | • | • | • | • |
| No master server overhead                       | • | • | • | • | • | • | • |
| No waiting for                                  | • | with sync off | with sync off | • | • | • | • |
### Feature

<table>
<thead>
<tr>
<th>Feature</th>
<th>Shared Disk Failover</th>
<th>File System Replication</th>
<th>Write-Ahead Log Shipping</th>
<th>Logical Replication</th>
<th>Trigger-Based Master-Standby Replication</th>
<th>Statement-Based Replication Middleware</th>
<th>Asynchronous Multi-master Replication</th>
<th>Synchronous Multi-master Replication</th>
</tr>
</thead>
<tbody>
<tr>
<td>multiple servers</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Master failure will never lose data</td>
<td>•</td>
<td>•</td>
<td>with sync on</td>
<td>with sync on</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Replicas accept read-only queries</td>
<td></td>
<td></td>
<td></td>
<td>with hot standby</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Per-table granularity</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>No conflict resolution necessary</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

There are a few solutions that do not fit into the above categories:

### Data Partitioning

Data partitioning splits tables into data sets. Each set can be modified by only one server. For example, data can be partitioned by offices, e.g., London and Paris, with a server in each office. If queries combining London and Paris data are necessary, an application can query both servers, or master/standby replication can be used to keep a read-only copy of the other office’s data on each server.

### Multiple-Server Parallel Query Execution

Many of the above solutions allow multiple servers to handle multiple queries, but none allow a single query to use multiple servers to complete faster. This solution allows multiple servers to work concurrently on a single query. It is usually accomplished by splitting the data among servers and having each server execute its part of the query and return results to a central server where they are combined and returned to the user. Pgpool-II has this capability. Also, this can be implemented using the PL/Proxy tool set.

## 26.2. Log-Shipping Standby Servers

Continuous archiving can be used to create a high availability (HA) cluster configuration with one or more standby servers ready to take over operations if the primary server fails. This capability is widely referred to as warm standby or log shipping.

The primary and standby server work together to provide this capability, though the servers are only loosely coupled. The primary server operates in continuous archiving mode, while each standby server operates in continuous recovery mode, reading the WAL files from the primary. No changes to the database tables are required to enable this capability, so it offers low administration overhead compared to some other replication solutions. This configuration also has relatively low performance impact on the primary server.

Directly moving WAL records from one database server to another is typically described as log shipping. PostgreSQL implements file-based log shipping by transferring WAL records one file (WAL...
High Availability, Load Balancing, and Replication

segment) at a time. WAL files (16MB) can be shipped easily and cheaply over any distance, whether it be to an adjacent system, another system at the same site, or another system on the far side of the globe. The bandwidth required for this technique varies according to the transaction rate of the primary server. Record-based log shipping is more granular and streams WAL changes incrementally over a network connection (see Section 26.2.5).

It should be noted that log shipping is asynchronous, i.e., the WAL records are shipped after transaction commit. As a result, there is a window for data loss should the primary server suffer a catastrophic failure; transactions not yet shipped will be lost. The size of the data loss window in file-based log shipping can be limited by use of the `archive_timeout` parameter, which can be set as low as a few seconds. However such a low setting will substantially increase the bandwidth required for file shipping. Streaming replication (see Section 26.2.5) allows a much smaller window of data loss.

Recovery performance is sufficiently good that the standby will typically be only moments away from full availability once it has been activated. As a result, this is called a warm standby configuration which offers high availability. Restoring a server from an archived base backup and rollforward will take considerably longer, so that technique only offers a solution for disaster recovery, not high availability. A standby server can also be used for read-only queries, in which case it is called a Hot Standby server. See Section 26.5 for more information.

26.2.1. Planning

It is usually wise to create the primary and standby servers so that they are as similar as possible, at least from the perspective of the database server. In particular, the path names associated with tablespaces will be passed across unmodified, so both primary and standby servers must have the same mount paths for tablespaces if that feature is used. Keep in mind that if `CREATE TABLESPACE` is executed on the primary, any new mount point needed for it must be created on the primary and all standby servers before the command is executed. Hardware need not be exactly the same, but experience shows that maintaining two identical systems is easier than maintaining two dissimilar ones over the lifetime of the application and system. In any case the hardware architecture must be the same — shipping from, say, a 32-bit to a 64-bit system will not work.

In general, log shipping between servers running different major PostgreSQL release levels is not possible. It is the policy of the PostgreSQL Global Development Group not to make changes to disk formats during minor release upgrades, so it is likely that running different minor release levels on primary and standby servers will work successfully. However, no formal support for that is offered and you are advised to keep primary and standby servers at the same release level as much as possible. When updating to a new minor release, the safest policy is to update the standby servers first — a new minor release is more likely to be able to read WAL files from a previous minor release than vice versa.

26.2.2. Standby Server Operation

In standby mode, the server continuously applies WAL received from the master server. The standby server can read WAL from a WAL archive (see `restore_command`) or directly from the master over a TCP connection (streaming replication). The standby server will also attempt to restore any WAL found in the standby cluster’s `pg_wal` directory. That typically happens after a server restart, when the standby replays again WAL that was streamed from the master before the restart, but you can also manually copy files to `pg_wal` at any time to have them replayed.

At startup, the standby begins by restoring all WAL available in the archive location, calling `restore_command`. Once it reaches the end of WAL available there and `restore_command` fails, it tries to restore any WAL available in the `pg_wal` directory. If that fails, and streaming replication has been configured, the standby tries to connect to the primary server and start streaming WAL from the last valid record found in archive or `pg_wal`. If that fails or streaming replication is not configured, or if the connection is later disconnected, the standby goes back to step 1 and tries to restore the file from the archive again. This loop of retries from the archive, `pg_wal`, and via streaming replication goes on until the server is stopped or failover is triggered by a trigger file.
Standby mode is exited and the server switches to normal operation when `pg_ctl promote` is run or a trigger file is found (`trigger_file`). Before failover, any WAL immediately available in the archive or in `pg_wal` will be restored, but no attempt is made to connect to the master.

### 26.2.3. Preparing the Master for Standby Servers

Set up continuous archiving on the primary to an archive directory accessible from the standby, as described in Section 25.3. The archive location should be accessible from the standby even when the master is down, i.e. it should reside on the standby server itself or another trusted server, not on the master server.

If you want to use streaming replication, set up authentication on the primary server to allow replication connections from the standby server(s); that is, create a role and provide a suitable entry or entries in `pg_hba.conf` with the database field set to `replication`. Also ensure `max_wal_senders` is set to a sufficiently large value in the configuration file of the primary server. If replication slots will be used, ensure that `max_replication_slots` is set sufficiently high as well.

Take a base backup as described in Section 25.3.2 to bootstrap the standby server.

### 26.2.4. Setting Up a Standby Server

To set up the standby server, restore the base backup taken from primary server (see Section 25.3.4). Create a recovery command file `recovery.conf` in the standby's cluster data directory, and turn on `standby_mode`. Set `restore_command` to a simple command to copy files from the WAL archive. If you plan to have multiple standby servers for high availability purposes, set `recovery_target_timeline` to `latest`, to make the standby server follow the timeline change that occurs at failover to another standby.

---

**Note**

Do not use `pg_standby` or similar tools with the built-in standby mode described here. `restore_command` should return immediately if the file does not exist; the server will retry the command again if necessary. See Section 26.4 for using tools like `pg_standby`.

---

If you want to use streaming replication, fill in `primary_conninfo` with a `libpq` connection string, including the host name (or IP address) and any additional details needed to connect to the primary server. If the primary needs a password for authentication, the password needs to be specified in `primary_conninfo` as well.

If you're setting up the standby server for high availability purposes, set up WAL archiving, connections and authentication like the primary server, because the standby server will work as a primary server after failover.

If you're using a WAL archive, its size can be minimized using the `archive_cleanup_command` parameter to remove files that are no longer required by the standby server. The `pg_archivecleanup` utility is designed specifically to be used with `archive_cleanup_command` in typical single-standby configurations, see `pg_archivecleanup`. Note however, that if you're using the archive for backup purposes, you need to retain files needed to recover from at least the latest base backup, even if they're no longer needed by the standby.

A simple example of a `recovery.conf` is:

```plaintext
standby_mode = 'on'
```
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo password=foopass'
restore_command = 'cp /path/to/archive/%f %p'
archive_cleanup_command = 'pg_archivecleanup /path/to/archive %r'

You can have any number of standby servers, but if you use streaming replication, make sure you set max_wal_senders high enough in the primary to allow them to be connected simultaneously.

### 26.2.5. Streaming Replication

Streaming replication allows a standby server to stay more up-to-date than is possible with file-based log shipping. The standby connects to the primary, which streams WAL records to the standby as they're generated, without waiting for the WAL file to be filled.

Streaming replication is asynchronous by default (see Section 26.2.8), in which case there is a small delay between committing a transaction in the primary and the changes becoming visible in the standby. This delay is however much smaller than with file-based log shipping, typically under one second assuming the standby is powerful enough to keep up with the load. With streaming replication, archive_timeout is not required to reduce the data loss window.

If you use streaming replication without file-based continuous archiving, the server might recycle old WAL segments before the standby has received them. If this occurs, the standby will need to be reinitialized from a new base backup. You can avoid this by setting wal_keep_segments to a value large enough to ensure that WAL segments are not recycled too early, or by configuring a replication slot for the standby. If you set up a WAL archive that's accessible from the standby, these solutions are not required, since the standby can always use the archive to catch up provided it retains enough segments.

To use streaming replication, set up a file-based log-shipping standby server as described in Section 26.2. The step that turns a file-based log-shipping standby into streaming replication standby is setting primary_conninfo setting in the recovery.conf file to point to the primary server. Set listen_addresses and authentication options (see pg_hba.conf) on the primary so that the standby server can connect to the replication pseudo-database on the primary server (see Section 26.2.5.1).

On systems that support the keepalive socket option, setting tcp_keepalives_idle, tcp_keepalives_interval and tcp_keepalives_count helps the primary promptly notice a broken connection.

Set the maximum number of concurrent connections from the standby servers (see max_wal_senders for details).

When the standby is started and primary_conninfo is set correctly, the standby will connect to the primary after replaying all WAL files available in the archive. If the connection is established successfully, you will see a walreceiver process in the standby, and a corresponding walsender process in the primary.

#### 26.2.5.1. Authentication

It is very important that the access privileges for replication be set up so that only trusted users can read the WAL stream, because it is easy to extract privileged information from it. Standby servers must authenticate to the primary as a superuser or an account that has the REPLICATION privilege. It is recommended to create a dedicated user account with REPLICATION and LOGIN privileges for replication. While REPLICATION privilege gives very high permissions, it does not allow the user to modify any data on the primary system, which the SUPERUSER privilege does.

Client authentication for replication is controlled by a pg_hba.conf record specifying replication in the database field. For example, if the standby is running on host IP 192.168.1.100 and the account name for replication is foo, the administrator can add the following line to the pg_hba.conf file on the primary:

```bash
host    all         all  trust
192.168.1.100 foo    replication
```
# Allow the user "foo" from host 192.168.1.100 to connect to the primary as a replication standby if the user's password is correctly supplied.

<table>
<thead>
<tr>
<th>TYPE</th>
<th>DATABASE</th>
<th>USER</th>
<th>ADDRESS</th>
<th>METHOD</th>
</tr>
</thead>
<tbody>
<tr>
<td>host</td>
<td>replication</td>
<td>foo</td>
<td>192.168.1.100/32</td>
<td>md5</td>
</tr>
</tbody>
</table>

The host name and port number of the primary, connection user name, and password are specified in the `recovery.conf` file. The password can also be set in the `~/.pgpass` file on the standby (specify replication in the `database` field). For example, if the primary is running on host IP 192.168.1.50, port 5432, the account name for replication is `foo`, and the password is `foopass`, the administrator can add the following line to the `recovery.conf` file on the standby:

```sql
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo password=foopass'
```

### 26.2.5.2. Monitoring

An important health indicator of streaming replication is the amount of WAL records generated in the primary, but not yet applied in the standby. You can calculate this lag by comparing the current WAL write location on the primary with the last WAL location received by the standby. These locations can be retrieved using `pg_current_wal_lsn` on the primary and `pg_last_wal_receive_lsn` on the standby, respectively (see Table 9.79 and Table 9.80 for details). The last WAL receive location in the standby is also displayed in the process status of the WAL receiver process, displayed using the `ps` command (see Section 28.1 for details).

You can retrieve a list of WAL sender processes via the `pg_stat_replication` view. Large differences between `pg_current_wal_lsn` and the view's `sent_lsn` field might indicate that the master server is under heavy load, while differences between `sent_lsn` and `pg_last_wal_receive_lsn` on the standby might indicate network delay, or that the standby is under heavy load.

On a hot standby, the status of the WAL receiver process can be retrieved via the `pg_stat_wal_receiver` view. A large difference between `pg_last_wal_replay_lsn` and the view's `received_lsn` indicates that WAL is being received faster than it can be replayed.

### 26.2.6. Replication Slots

Replication slots provide an automated way to ensure that the master does not remove WAL segments until they have been received by all standbys, and that the master does not remove rows which could cause a recovery conflict even when the standby is disconnected.

In lieu of using replication slots, it is possible to prevent the removal of old WAL segments using `wal_keep_segments`, or by storing the segments in an archive using `archive_command`. However, these methods often result in retaining more WAL segments than required, whereas replication slots retain only the number of segments known to be needed. An advantage of these methods is that they bound the space requirement for `pg_wal`; there is currently no way to do this using replication slots.

Similarly, `hot_standby_feedback` and `vacuum_defer_cleanup_age` provide protection against relevant rows being removed by vacuum, but the former provides no protection during any time period when the standby is not connected, and the latter often needs to be set to a high value to provide adequate protection. Replication slots overcome these disadvantages.
26.2.6.1. Querying and manipulating replication slots

Each replication slot has a name, which can contain lower-case letters, numbers, and the underscore character.

Existing replication slots and their state can be seen in the `pg_replication_slots` view.

Slots can be created and dropped either via the streaming replication protocol (see Section 53.4) or via SQL functions (see Section 9.26.6).

26.2.6.2. Configuration Example

You can create a replication slot like this:

```
postgres=# SELECT * FROM
    pg_create_physical_replication_slot('node_a_slot');
```

```
<table>
<thead>
<tr>
<th>slot_name</th>
<th>lsn</th>
</tr>
</thead>
<tbody>
<tr>
<td>node_a_slot</td>
<td>-</td>
</tr>
</tbody>
</table>
```

```
postgres=# SELECT slot_name, slot_type, active FROM
    pg_replication_slots;
```

```
<table>
<thead>
<tr>
<th>slot_name</th>
<th>slot_type</th>
<th>active</th>
</tr>
</thead>
<tbody>
<tr>
<td>node_a_slot</td>
<td>physical</td>
<td>f</td>
</tr>
</tbody>
</table>
```

To configure the standby to use this slot, `primary_slot_name` should be configured in the standby's `recovery.conf`. Here is a simple example:

```
standby_mode = 'on'
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo
    password=foopass'
primary_slot_name = 'node_a_slot'
```

26.2.7. Cascading Replication

The cascading replication feature allows a standby server to accept replication connections and stream WAL records to other standbys, acting as a relay. This can be used to reduce the number of direct connections to the master and also to minimize inter-site bandwidth overheads.

A standby acting as both a receiver and a sender is known as a cascading standby. Standbys that are more directly connected to the master are known as upstream servers, while those standby servers further away are downstream servers. Cascading replication does not place limits on the number or arrangement of downstream servers, though each standby connects to only one upstream server which eventually links to a single master/primary server.

A cascading standby sends not only WAL records received from the master but also those restored from the archive. So even if the replication connection in some upstream connection is terminated, streaming replication continues downstream for as long as new WAL records are available.

Cascading replication is currently asynchronous. Synchronous replication (see Section 26.2.8) settings have no effect on cascading replication at present.

Hot Standby feedback propagates upstream, whatever the cascaded arrangement.
If an upstream standby server is promoted to become new master, downstream servers will continue to stream from the new master if `recovery_target_timeline` is set to 'latest'.

To use cascading replication, set up the cascading standby so that it can accept replication connections (that is, set `max_wal_senders` and `hot_standby`, and configure host-based authentication). You will also need to set `primary_conninfo` in the downstream standby to point to the cascading standby.

### 26.2.8. Synchronous Replication

PostgreSQL streaming replication is asynchronous by default. If the primary server crashes then some transactions that were committed may not have been replicated to the standby server, causing data loss. The amount of data loss is proportional to the replication delay at the time of failover.

Synchronous replication offers the ability to confirm that all changes made by a transaction have been transferred to one or more synchronous standby servers. This extends that standard level of durability offered by a transaction commit. This level of protection is referred to as 2-safe replication in computer science theory, and group-1-safe (group-safe and 1-safe) when `synchronous_commit` is set to `remote_write`.

When requesting synchronous replication, each commit of a write transaction will wait until confirmation is received that the commit has been written to the write-ahead log on disk of both the primary and standby server. The only possibility that data can be lost is if both the primary and the standby suffer crashes at the same time. This can provide a much higher level of durability, though only if the sysadmin is cautious about the placement and management of the two servers. Waiting for confirmation increases the user's confidence that the changes will not be lost in the event of server crashes but it also necessarily increases the response time for the requesting transaction. The minimum wait time is the round-trip time between primary to standby.

Read only transactions and transaction rollbacks need not wait for replies from standby servers. Sub-transaction commits do not wait for responses from standby servers, only top-level commits. Long running actions such as data loading or index building do not wait until the very final commit message. All two-phase commit actions require commit waits, including both prepare and commit.

A synchronous standby can be a physical replication standby or a logical replication subscriber. It can also be any other physical or logical WAL replication stream consumer that knows how to send the appropriate feedback messages. Besides the built-in physical and logical replication systems, this includes special programs such as `pg_receivewal` and `pg_recvlogical` as well as some third-party replication systems and custom programs. Check the respective documentation for details on synchronous replication support.

#### 26.2.8.1. Basic Configuration

Once streaming replication has been configured, configuring synchronous replication requires only one additional configuration step: `synchronous_standby_names` must be set to a non-empty value. `synchronous_commit` must also be set to `on`, but since this is the default value, typically no change is required. (See Section 19.5.1 and Section 19.6.2.) This configuration will cause each commit to wait for confirmation that the standby has written the commit record to durable storage. `synchronous_commit` can be set by individual users, so it can be configured in the configuration file, for particular users or databases, or dynamically by applications, in order to control the durability guarantee on a per-transaction basis.

After a commit record has been written to disk on the primary, the WAL record is then sent to the standby. The standby sends reply messages each time a new batch of WAL data is written to disk, unless `wal_receiver_status_interval` is set to zero on the standby. In the case that `synchronous_commit` is set to `remote_apply`, the standby sends reply messages when the commit record is replayed, making the transaction visible. If the standby is chosen as a synchronous standby, according to the setting of `synchronous_standby_names` on the primary, the reply messages from that standby will be considered along with those from other synchronous standbys to decide when to release transactions waiting for confirmation that the commit record has been received. These
parameters allow the administrator to specify which standby servers should be synchronous standbys. Note that the configuration of synchronous replication is mainly on the master. Named standbys must be directly connected to the master; the master knows nothing about downstream standby servers using cascaded replication.

Setting `synchronous_commit` to `remote_write` will cause each commit to wait for confirmation that the standby has received the commit record and written it out to its own operating system, but not for the data to be flushed to disk on the standby. This setting provides a weaker guarantee of durability than `on` does; the standby could lose the data in the event of an operating system crash, though not a PostgreSQL crash. However, it’s a useful setting in practice because it can decrease the response time for the transaction. Data loss could only occur if both the primary and the standby crash and the database of the primary gets corrupted at the same time.

Setting `synchronous_commit` to `remote_apply` will cause each commit to wait until the current synchronous standbys report that they have replayed the transaction, making it visible to user queries. In simple cases, this allows for load balancing with causal consistency.

Users will stop waiting if a fast shutdown is requested. However, as when using asynchronous replication, the server will not fully shutdown until all outstanding WAL records are transferred to the currently connected standby servers.

### 26.2.8.2. Multiple Synchronous Standbys

Synchronous replication supports one or more synchronous standby servers; transactions will wait until all the standby servers which are considered as synchronous confirm receipt of their data. The number of synchronous standbys that transactions must wait for replies from is specified in `synchronous_standby_names`. This parameter also specifies a list of standby names and the method (FIRST and ANY) to choose synchronous standbys from the listed ones.

The method `FIRST` specifies a priority-based synchronous replication and makes transaction commits wait until their WAL records are replicated to the requested number of synchronous standbys chosen based on their priorities. The standbys whose names appear earlier in the list are given higher priority and will be considered as synchronous. Other standby servers appearing later in this list represent potential synchronous standbys. If any of the current synchronous standbys disconnects for whatever reason, it will be replaced immediately with the next-highest-priority standby.

An example of `synchronous_standby_names` for a priority-based multiple synchronous standbys is:

```
synchronous_standby_names = 'FIRST 2 (s1, s2, s3)'
```

In this example, if four standby servers `s1`, `s2`, `s3` and `s4` are running, the two standbys `s1` and `s2` will be chosen as synchronous standbys because their names appear early in the list of standby names. `s3` is a potential synchronous standby and will take over the role of synchronous standby when either of `s1` or `s2` fails. `s4` is an asynchronous standby since its name is not in the list.

The method `ANY` specifies a quorum-based synchronous replication and makes transaction commits wait until their WAL records are replicated to at least the requested number of synchronous standbys in the list.

An example of `synchronous_standby_names` for a quorum-based multiple synchronous standbys is:

```
synchronous_standby_names = 'ANY 2 (s1, s2, s3)'
```

In this example, if four standby servers `s1`, `s2`, `s3` and `s4` are running, transaction commits will wait for replies from at least any two standbys of `s1`, `s2` and `s3`. `s4` is an asynchronous standby since its name is not in the list.
The synchronous states of standby servers can be viewed using the `pg_stat_replication` view.

## 26.2.8.3. Planning for Performance

Synchronous replication usually requires carefully planned and placed standby servers to ensure applications perform acceptably. Waiting doesn't utilize system resources, but transaction locks continue to be held until the transfer is confirmed. As a result, incautious use of synchronous replication will reduce performance for database applications because of increased response times and higher contention.

PostgreSQL allows the application developer to specify the durability level required via replication. This can be specified for the system overall, though it can also be specified for specific users or connections, or even individual transactions.

For example, an application workload might consist of: 10% of changes are important customer details, while 90% of changes are less important data that the business can more easily survive if it is lost, such as chat messages between users.

With synchronous replication options specified at the application level (on the primary) we can offer synchronous replication for the most important changes, without slowing down the bulk of the total workload. Application level options are an important and practical tool for allowing the benefits of synchronous replication for high performance applications.

You should consider that the network bandwidth must be higher than the rate of generation of WAL data.

## 26.2.8.4. Planning for High Availability

`synchronous_standby_names` specifies the number and names of synchronous standbys that transaction commits made when `synchronous_commit` is set to `on`, `remote_apply` or `remote_write` will wait for responses from. Such transaction commits may never be completed if any one of synchronous standbys should crash.

The best solution for high availability is to ensure you keep as many synchronous standbys as requested. This can be achieved by naming multiple potential synchronous standbys using `synchronous_standby_names`.

In a priority-based synchronous replication, the standbys whose names appear earlier in the list will be used as synchronous standbys. Standbys listed after these will take over the role of synchronous standby if one of current ones should fail.

In a quorum-based synchronous replication, all the standbys appearing in the list will be used as candidates for synchronous standbys. Even if one of them should fail, the other standbys will keep performing the role of candidates of synchronous standby.

When a standby first attaches to the primary, it will not yet be properly synchronized. This is described as `catchup` mode. Once the lag between standby and primary reaches zero for the first time we move to real-time `streaming` state. The catch-up duration may be long immediately after the standby has been created. If the standby is shut down, then the catch-up period will increase according to the length of time the standby has been down. The standby is only able to become a synchronous standby once it has reached `streaming` state. This state can be viewed using the `pg_stat_replication` view.

If primary restarts while commits are waiting for acknowledgement, those waiting transactions will be marked fully committed once the primary database recovers. There is no way to be certain that all standbys have received all outstanding WAL data at time of the crash of the primary. Some transactions may not show as committed on the standby, even though they show as committed on the primary. The guarantee we offer is that the application will not receive explicit acknowledgement of the successful commit of a transaction until the WAL data is known to be safely received by all the synchronous standbys.

679
If you really cannot keep as many synchronous standbys as requested then you should decrease the number of synchronous standbys that transaction commits must wait for responses from in `synchronous_standby_names` (or disable it) and reload the configuration file on the primary server.

If the primary is isolated from remaining standby servers you should fail over to the best candidate of those other remaining standby servers.

If you need to re-create a standby server while transactions are waiting, make sure that the commands `pg_start_backup()` and `pg_stop_backup()` are run in a session with `synchronous_commit = off`, otherwise those requests will wait forever for the standby to appear.

### 26.2.9. Continuous archiving in standby

When continuous WAL archiving is used in a standby, there are two different scenarios: the WAL archive can be shared between the primary and the standby, or the standby can have its own WAL archive. When the standby has its own WAL archive, set `archive_mode` to `always`, and the standby will call the archive command for every WAL segment it receives, whether it's by restoring from the archive or by streaming replication. The shared archive can be handled similarly, but the `archive_command` must test if the file being archived exists already, and if the existing file has identical contents. This requires more care in the `archive_command`, as it must be careful to not overwrite an existing file with different contents, but return success if the exactly same file is archived twice. And all that must be done free of race conditions, if two servers attempt to archive the same file at the same time.

If `archive_mode` is set to `on`, the archiver is not enabled during recovery or standby mode. If the standby server is promoted, it will start archiving after the promotion, but will not archive any WAL it did not generate itself. To get a complete series of WAL files in the archive, you must ensure that all WAL is archived, before it reaches the standby. This is inherently true with file-based log shipping, as the standby can only restore files that are found in the archive, but not if streaming replication is enabled. When a server is not in recovery mode, there is no difference between `on` and `always` modes.

### 26.3. Failover

If the primary server fails then the standby server should begin failover procedures.

If the standby server fails then no failover need take place. If the standby server can be restarted, even some time later, then the recovery process can also be restarted immediately, taking advantage of restartable recovery. If the standby server cannot be restarted, then a full new standby server instance should be created.

If the primary server fails and the standby server becomes the new primary, and then the old primary restarts, you must have a mechanism for informing the old primary that it is no longer the primary. This is sometimes known as STONITH (Shoot The Other Node In The Head), which is necessary to avoid situations where both systems think they are the primary, which will lead to confusion and ultimately data loss.

Many failover systems use just two systems, the primary and the standby, connected by some kind of heartbeat mechanism to continually verify the connectivity between the two and the viability of the primary. It is also possible to use a third system (called a witness server) to prevent some cases of inappropriate failover, but the additional complexity might not be worthwhile unless it is set up with sufficient care and rigorous testing.

PostgreSQL does not provide the system software required to identify a failure on the primary and notify the standby database server. Many such tools exist and are well integrated with the operating system facilities required for successful failover, such as IP address migration.

Once failover to the standby occurs, there is only a single server in operation. This is known as a degenerate state. The former standby is now the primary, but the former primary is down and might
stay down. To return to normal operation, a standby server must be recreated, either on the former primary system when it comes up, or on a third, possibly new, system. The pg_rewind utility can be used to speed up this process on large clusters. Once complete, the primary and standby can be considered to have switched roles. Some people choose to use a third server to provide backup for the new primary until the new standby server is recreated, though clearly this complicates the system configuration and operational processes.

So, switching from primary to standby server can be fast but requires some time to re-prepare the failover cluster. Regular switching from primary to standby is useful, since it allows regular downtime on each system for maintenance. This also serves as a test of the failover mechanism to ensure that it will really work when you need it. Written administration procedures are advised.

To trigger failover of a log-shipping standby server, run pg_ctl promote or create a trigger file with the file name and path specified by the trigger_file setting in recovery.conf. If you're planning to use pg_ctl promote to fail over, trigger_file is not required. If you're setting up the reporting servers that are only used to offload read-only queries from the primary, not for high availability purposes, you don’t need to promote it.

### 26.4. Alternative Method for Log Shipping

An alternative to the built-in standby mode described in the previous sections is to use a restore_command that polls the archive location. This was the only option available in versions 8.4 and below. In this setup, set standby_mode off, because you are implementing the polling required for standby operation yourself. See the pg_standby module for a reference implementation of this.

Note that in this mode, the server will apply WAL one file at a time, so if you use the standby server for queries (see Hot Standby), there is a delay between an action in the master and when the action becomes visible in the standby, corresponding the time it takes to fill up the WAL file. archive_timeout can be used to make that delay shorter. Also note that you can’t combine streaming replication with this method.

The operations that occur on both primary and standby servers are normal continuous archiving and recovery tasks. The only point of contact between the two database servers is the archive of WAL files that both share: primary writing to the archive, standby reading from the archive. Care must be taken to ensure that WAL archives from separate primary servers do not become mixed together or confused. The archive need not be large if it is only required for standby operation.

The magic that makes the two loosely coupled servers work together is simply a restore_command used on the standby that, when asked for the next WAL file, waits for it to become available from the primary. The restore_command is specified in the recovery.conf file on the standby server. Normal recovery processing would request a file from the WAL archive, reporting failure if the file was unavailable. For standby processing it is normal for the next WAL file to be unavailable, so the standby must wait for it to appear. For files ending in .history there is no need to wait, and a non-zero return code must be returned. A waiting restore_command can be written as a custom script that loops after polling for the existence of the next WAL file. There must also be some way to trigger failover, which should interrupt the restore_command, break the loop and return a file-not-found error to the standby server. This ends recovery and the standby will then come up as a normal server.

Pseudocode for a suitable restore_command is:

```c
triggered = false;
while (!NextWALFileReady() && !triggered)
{
  sleep(100000L); /* wait for ~0.1 sec */
  if (CheckForExternalTrigger())
    triggered = true;
}
if (!triggered)
```
CopyWALFileForRecovery();

A working example of a waiting restore_command is provided in the pg_standby module. It should be used as a reference on how to correctly implement the logic described above. It can also be extended as needed to support specific configurations and environments.

The method for triggering failover is an important part of planning and design. One potential option is the restore_command command. It is executed once for each WAL file, but the process running the restore_command is created and dies for each file, so there is no daemon or server process, and signals or a signal handler cannot be used. Therefore, the restore_command is not suitable to trigger failover. It is possible to use a simple timeout facility, especially if used in conjunction with a known archive_timeout setting on the primary. However, this is somewhat error prone since a network problem or busy primary server might be sufficient to initiate failover. A notification mechanism such as the explicit creation of a trigger file is ideal, if this can be arranged.

26.4.1. Implementation

The short procedure for configuring a standby server using this alternative method is as follows. For full details of each step, refer to previous sections as noted.

1. Set up primary and standby systems as nearly identical as possible, including two identical copies of PostgreSQL at the same release level.

2. Set up continuous archiving from the primary to a WAL archive directory on the standby server. Ensure that archive_mode, archive_command and archive_timeout are set appropriately on the primary (see Section 25.3.1).

3. Make a base backup of the primary server (see Section 25.3.2), and load this data onto the standby.

4. Begin recovery on the standby server from the local WAL archive, using a recovery.conf that specifies a restore_command that waits as described previously (see Section 25.3.4).

Recovery treats the WAL archive as read-only, so once a WAL file has been copied to the standby system it can be copied to tape at the same time as it is being read by the standby database server. Thus, running a standby server for high availability can be performed at the same time as files are stored for longer term disaster recovery purposes.

For testing purposes, it is possible to run both primary and standby servers on the same system. This does not provide any worthwhile improvement in server robustness, nor would it be described as HA.

26.4.2. Record-based Log Shipping

It is also possible to implement record-based log shipping using this alternative method, though this requires custom development, and changes will still only become visible to hot standby queries after a full WAL file has been shipped.

An external program can call the pg_walfile_name_offset() function (see Section 9.26) to find out the file name and the exact byte offset within it of the current end of WAL. It can then access the WAL file directly and copy the data from the last known end of WAL through the current end over to the standby servers. With this approach, the window for data loss is the polling cycle time of the copying program, which can be very small, and there is no wasted bandwidth from forcing partially-used segment files to be archived. Note that the standby servers' restore_command scripts can only deal with whole WAL files, so the incrementally copied data is not ordinarily made available to the standby servers. It is of use only when the primary dies — then the last partial WAL file is fed to the standby before allowing it to come up. The correct implementation of this process requires cooperation of the restore_command script with the data copying program.

Starting with PostgreSQL version 9.0, you can use streaming replication (see Section 26.2.5) to achieve the same benefits with less effort.
26.5. Hot Standby

Hot Standby is the term used to describe the ability to connect to the server and run read-only queries while the server is in archive recovery or standby mode. This is useful both for replication purposes and for restoring a backup to a desired state with great precision. The term Hot Standby also refers to the ability of the server to move from recovery through to normal operation while users continue running queries and/or keep their connections open.

Running queries in hot standby mode is similar to normal query operation, though there are several usage and administrative differences explained below.

26.5.1. User's Overview

When the hot_standby parameter is set to true on a standby server, it will begin accepting connections once the recovery has brought the system to a consistent state. All such connections are strictly read-only; not even temporary tables may be written.

The data on the standby takes some time to arrive from the primary server so there will be a measurable delay between primary and standby. Running the same query nearly simultaneously on both primary and standby might therefore return differing results. We say that data on the standby is eventually consistent with the primary. Once the commit record for a transaction is replayed on the standby, the changes made by that transaction will be visible to any new snapshots taken on the standby. Snapshots may be taken at the start of each query or at the start of each transaction, depending on the current transaction isolation level. For more details, see Section 13.2.

Transactions started during hot standby may issue the following commands:

- Query access - SELECT, COPY TO
- Cursor commands - DECLARE, FETCH, CLOSE
- Parameters - SHOW, SET, RESET
- Transaction management commands
  - BEGIN, END, ABORT, START TRANSACTION
  - SAVEPOINT, RELEASE, ROLLBACK TO SAVEPOINT
  - EXCEPTION blocks and other internal subtransactions
  - LOCK TABLE, though only when explicitly in one of these modes: ACCESS SHARE, ROW SHARE or ROW EXCLUSIVE.
- Plans and resources - PREPARE, EXECUTE, DEALLOCATE, DISCARD
- Plugins and extensions - LOAD

Transactions started during hot standby will never be assigned a transaction ID and cannot write to the system write-ahead log. Therefore, the following actions will produce error messages:

- Data Manipulation Language (DML) - INSERT, UPDATE, DELETE, COPY FROM, TRUNCATE. Note that there are no allowed actions that result in a trigger being executed during recovery. This restriction applies even to temporary tables, because table rows cannot be read or written without assigning a transaction ID, which is currently not possible in a Hot Standby environment.
- Data Definition Language (DDL) - CREATE, DROP, ALTER, COMMENT. This restriction applies even to temporary tables, because carrying out these operations would require updating the system catalog tables.
SELECT ... FOR SHARE | UPDATE, because row locks cannot be taken without updating the underlying data files.

Rules on SELECT statements that generate DML commands.

LOCK that explicitly requests a mode higher than ROW EXCLUSIVE MODE.

LOCK in short default form, since it requests ACCESS EXCLUSIVE MODE.

Transaction management commands that explicitly set non-read-only state:

BEGIN READ WRITE, START TRANSACTION READ WRITE

SET TRANSACTION READ WRITE, SET SESSION CHARACTERISTICS AS TRANSACTION READ WRITE

SET transaction_read_only = off

Two-phase commit commands - PREPARE TRANSACTION, COMMIT PREPARED, ROLLBACK PREPARED because even read-only transactions need to write WAL in the prepare phase (the first phase of two phase commit).

Sequence updates - nextval(), setval()

LISTEN, UNLISTEN, NOTIFY

In normal operation, “read-only” transactions are allowed to use LISTEN, UNLISTEN, and NOTIFY, so Hot Standby sessions operate under slightly tighter restrictions than ordinary read-only sessions. It is possible that some of these restrictions might be loosened in a future release.

During hot standby, the parameter transaction_read_only is always true and may not be changed. But as long as no attempt is made to modify the database, connections during hot standby will act much like any other database connection. If failover or switchover occurs, the database will switch to normal processing mode. Sessions will remain connected while the server changes mode. Once hot standby finishes, it will be possible to initiate read-write transactions (even from a session begun during hot standby).

Users will be able to tell whether their session is read-only by issuing SHOW transaction_read_only. In addition, a set of functions (Table 9.80) allow users to access information about the standby server. These allow you to write programs that are aware of the current state of the database. These can be used to monitor the progress of recovery, or to allow you to write complex programs that restore the database to particular states.

26.5.2. Handling Query Conflicts

The primary and standby servers are in many ways loosely connected. Actions on the primary will have an effect on the standby. As a result, there is potential for negative interactions or conflicts between them. The easiest conflict to understand is performance: if a huge data load is taking place on the primary then this will generate a similar stream of WAL records on the standby, so standby queries may contend for system resources, such as I/O.

There are also additional types of conflict that can occur with Hot Standby. These conflicts are **hard conflicts** in the sense that queries might need to be canceled and, in some cases, sessions disconnected to resolve them. The user is provided with several ways to handle these conflicts. Conflict cases include:

- Access Exclusive locks taken on the primary server, including both explicit LOCK commands and various DDL actions, conflict with table accesses in standby queries.

- Dropping a tablespace on the primary conflicts with standby queries using that tablespace for temporary work files.
• Dropping a database on the primary conflicts with sessions connected to that database on the standby.

• Application of a vacuum cleanup record from WAL conflicts with standby transactions whose snapshots can still “see” any of the rows to be removed.

• Application of a vacuum cleanup record from WAL conflicts with queries accessing the target page on the standby, whether or not the data to be removed is visible.

On the primary server, these cases simply result in waiting; and the user might choose to cancel either of the conflicting actions. However, on the standby there is no choice: the WAL-logged action already occurred on the primary so the standby must not fail to apply it. Furthermore, allowing WAL application to wait indefinitely may be very undesirable, because the standby's state will become increasingly far behind the primary's. Therefore, a mechanism is provided to forcibly cancel standby queries that conflict with to-be-applied WAL records.

An example of the problem situation is an administrator on the primary server running `DROP TABLE` on a table that is currently being queried on the standby server. Clearly the standby query cannot continue if the `DROP TABLE` is applied on the standby. If this situation occurred on the primary, the `DROP TABLE` would wait until the other query had finished. But when `DROP TABLE` is run on the primary, the primary doesn't have information about what queries are running on the standby, so it will not wait for any such standby queries. The WAL change records come through to the standby while the standby query is still running, causing a conflict. The standby server must either delay application of the WAL records (and everything after them, too) or else cancel the conflicting query so that the `DROP TABLE` can be applied.

When a conflicting query is short, it's typically desirable to allow it to complete by delaying WAL application for a little bit; but a long delay in WAL application is usually not desirable. So the cancel mechanism has parameters, `max_standby_archive_delay` and `max_standby_streaming_delay`, that define the maximum allowed delay in WAL application. Conflicting queries will be canceled once it has taken longer than the relevant delay setting to apply any newly-received WAL data. There are two parameters so that different delay values can be specified for the case of reading WAL data from an archive (i.e., initial recovery from a base backup or “catching up” a standby server that has fallen far behind) versus reading WAL data via streaming replication.

In a standby server that exists primarily for high availability, it's best to set the delay parameters relatively short, so that the server cannot fall far behind the primary due to delays caused by standby queries. However, if the standby server is meant for executing long-running queries, then a high or even infinite delay value may be preferable. Keep in mind however that a long-running query could cause other sessions on the standby server to not see recent changes on the primary, if it delays application of WAL records.

Once the delay specified by `max_standby_archive_delay` or `max_standby_streaming_delay` has been exceeded, conflicting queries will be canceled. This usually results just in a cancellation error, although in the case of replaying a `DROP DATABASE` the entire conflicting session will be terminated. Also, if the conflict is over a lock held by an idle transaction, the conflicting session is terminated (this behavior might change in the future).

Canceled queries may be retried immediately (after beginning a new transaction, of course). Since query cancellation depends on the nature of the WAL records being replayed, a query that was canceled may well succeed if it is executed again.

Keep in mind that the delay parameters are compared to the elapsed time since the WAL data was received by the standby server. Thus, the grace period allowed to any one query on the standby is never more than the delay parameter, and could be considerably less if the standby has already fallen behind as a result of waiting for previous queries to complete, or as a result of being unable to keep up with a heavy update load.

The most common reason for conflict between standby queries and WAL replay is “early cleanup”. Normally, PostgreSQL allows cleanup of old row versions when there are no transactions that need
to see them to ensure correct visibility of data according to MVCC rules. However, this rule can only be applied for transactions executing on the master. So it is possible that cleanup on the master will remove row versions that are still visible to a transaction on the standby.

Experienced users should note that both row version cleanup and row version freezing will potentially conflict with standby queries. Running a manual `VACUUM FREEZE` is likely to cause conflicts even on tables with no updated or deleted rows.

Users should be clear that tables that are regularly and heavily updated on the primary server will quickly cause cancellation of longer running queries on the standby. In such cases the setting of a finite value for `max_standby_archive_delay` or `max_standby_streaming_delay` can be considered similar to setting `statement_timeout`.

Remedial possibilities exist if the number of standby-query cancellations is found to be unacceptable. The first option is to set the parameter `hot_standby_feedback`, which prevents `VACUUM` from removing recently-dead rows and so cleanup conflicts do not occur. If you do this, you should note that this will delay cleanup of dead rows on the primary, which may result in undesirable table bloat. However, the cleanup situation will be no worse than if the standby queries were running directly on the primary server, and you are still getting the benefit of off-loading execution onto the standby. If standby servers connect and disconnect frequently, you might want to make adjustments to handle the period when `hot_standby_feedback` feedback is not being provided. For example, consider increasing `max_standby_archive_delay` so that queries are not rapidly canceled by conflicts in WAL archive files during disconnected periods. You should also consider increasing `max_standby_streaming_delay` to avoid rapid cancellations by newly-arrived streaming WAL entries after reconnection.

Another option is to increase `vacuum_defer_cleanup_age` on the primary server, so that dead rows will not be cleaned up as quickly as they normally would be. This will allow more time for queries to execute before they are canceled on the standby, without having to set a high `max_standby_streaming_delay`. However it is difficult to guarantee any specific execution-time window with this approach, since `vacuum_defer_cleanup_age` is measured in transactions executed on the primary server.

The number of query cancels and the reason for them can be viewed using the `pg_stat_database_conflicts` system view on the standby server. The `pg_stat_database` system view also contains summary information.

### 26.5.3. Administrator’s Overview

If `hot_standby` is on in `postgresql.conf` (the default value) and there is a `recovery.conf` file present, the server will run in Hot Standby mode. However, it may take some time for Hot Standby connections to be allowed, because the server will not accept connections until it has completed sufficient recovery to provide a consistent state against which queries can run. During this period, clients that attempt to connect will be refused with an error message. To confirm the server has come up, either loop trying to connect from the application, or look for these messages in the server logs:

```
LOG:  entering standby mode

... then some time later ...

LOG:  consistent recovery state reached
LOG:  database system is ready to accept read only connections
```

Consistency information is recorded once per checkpoint on the primary. It is not possible to enable hot standby when reading WAL written during a period when `wal_level` was not set to `replica` or `logical` on the primary. Reaching a consistent state can also be delayed in the presence of both of these conditions:

- A write transaction has more than 64 subtransactions
• Very long-lived write transactions

If you are running file-based log shipping ("warm standby"), you might need to wait until the next WAL file arrives, which could be as long as the `archive_timeout` setting on the primary.

The setting of some parameters on the standby will need reconfiguration if they have been changed on the primary. For these parameters, the value on the standby must be equal to or greater than the value on the primary. Therefore, if you want to increase these values, you should do so on all standby servers first, before applying the changes to the primary server. Conversely, if you want to decrease these values, you should do so on the primary server first, before applying the changes to all standby servers. If these parameters are not set high enough then the standby will refuse to start. Higher values can then be supplied and the server restarted to begin recovery again. These parameters are:

• `max_connections`
• `max_prepared_transactions`
• `max_locks_per_transaction`
• `max_worker_processes`

It is important that the administrator select appropriate settings for `max_standby_archive_delay` and `max_standby_streaming_delay`. The best choices vary depending on business priorities. For example if the server is primarily tasked as a High Availability server, then you will want low delay settings, perhaps even zero, though that is a very aggressive setting. If the standby server is tasked as an additional server for decision support queries then it might be acceptable to set the maximum delay values to many hours, or even -1 which means wait forever for queries to complete.

Transaction status "hint bits" written on the primary are not WAL-logged, so data on the standby will likely re-write the hints again on the standby. Thus, the standby server will still perform disk writes even though all users are read-only; no changes occur to the data values themselves. Users will still write large sort temporary files and re-generate relcache info files, so no part of the database is truly read-only during hot standby mode. Note also that writes to remote databases using dblink module, and other operations outside the database using PL functions will still be possible, even though the transaction is read-only locally.

The following types of administration commands are not accepted during recovery mode:

• Data Definition Language (DDL) - e.g. `CREATE INDEX`
• Privilege and Ownership - `GRANT`, `REVOKE`, `REASSIGN`
• Maintenance commands - `ANALYZE`, `VACUUM`, `CLUSTER`, `REINDEX`

Again, note that some of these commands are actually allowed during "read only" mode transactions on the primary.

As a result, you cannot create additional indexes that exist solely on the standby, nor statistics that exist solely on the standby. If these administration commands are needed, they should be executed on the primary, and eventually those changes will propagate to the standby.

`pg_cancel_backend()` and `pg_terminate_backend()` will work on user backends, but not the Startup process, which performs recovery. `pg_stat_activity` does not show recovering transactions as active. As a result, `pg_prepared_xacts` is always empty during recovery. If you wish to resolve in-doubt prepared transactions, view `pg_prepared_xacts` on the primary and issue commands to resolve transactions there or resolve them after the end of recovery.

`pg_locks` will show locks held by backends, as normal. `pg_locks` also shows a virtual transaction managed by the Startup process that owns all `AccessExclusiveLocks` held by transactions being replayed by recovery. Note that the Startup process does not acquire locks to make database changes,
and thus locks other than AccessExclusiveLocks do not show in pg_locks for the Startup process; they are just presumed to exist.

The Nagios plugin check_pgsqll will work, because the simple information it checks for exists. The check_postgres monitoring script will also work, though some reported values could give different or confusing results. For example, last vacuum time will not be maintained, since no vacuum occurs on the standby. Vacuums running on the primary do still send their changes to the standby.

WAL file control commands will not work during recovery, e.g. pg_start_backup, pg_switch_wal etc.

Dynamically loadable modules work, including pg_stat_statements.

Advisory locks work normally in recovery, including deadlock detection. Note that advisory locks are never WAL logged, so it is impossible for an advisory lock on either the primary or the standby to conflict with WAL replay. Nor is it possible to acquire an advisory lock on the primary and have it initiate a similar advisory lock on the standby. Advisory locks relate only to the server on which they are acquired.

Trigger-based replication systems such as Slony, Londiste and Bucardo won’t run on the standby at all, though they will run happily on the primary server as long as the changes are not sent to standby servers to be applied. WAL replay is not trigger-based so you cannot relay from the standby to any system that requires additional database writes or relies on the use of triggers.

New OIDs cannot be assigned, though some UUID generators may still work as long as they do not rely on writing new status to the database.

Currently, temporary table creation is not allowed during read only transactions, so in some cases existing scripts will not run correctly. This restriction might be relaxed in a later release. This is both a SQL Standard compliance issue and a technical issue.

DROP TABLESPACE can only succeed if the tablespace is empty. Some standby users may be actively using the tablespace via their temp_tablespaces parameter. If there are temporary files in the tablescape, all active queries are canceled to ensure that temporary files are removed, so the tablespace can be removed and WAL replay can continue.

Running DROP DATABASE or ALTER DATABASE ... SET TABLESPACE on the primary will generate a WAL entry that will cause all users connected to that database on the standby to be forcibly disconnected. This action occurs immediately, whatever the setting of max_standby_streaming_delay. Note that ALTER DATABASE ... RENAME does not disconnect users, which in most cases will go unnoticed, though might in some cases cause a program confusion if it depends in some way upon database name.

In normal (non-recovery) mode, if you issue DROP USER or DROP ROLE for a role with login capability while that user is still connected then nothing happens to the connected user - they remain connected. The user cannot reconnect however. This behavior applies in recovery also, so a DROP USER on the primary does not disconnect that user on the standby.

The statistics collector is active during recovery. All scans, reads, blocks, index usage, etc., will be recorded normally on the standby. Replayed actions will not duplicate their effects on primary, so replaying an insert will not increment the Inserts column of pg_stat_user_tables. The stats file is deleted at the start of recovery, so stats from primary and standby will differ; this is considered a feature, not a bug.

Autovacuum is not active during recovery. It will start normally at the end of recovery.

The background writer is active during recovery and will perform restartpoints (similar to checkpoints on the primary) and normal block cleaning activities. This can include updates of the hint bit information stored on the standby server. The CHECKPOINT command is accepted during recovery, though it performs a restartpoint rather than a new checkpoint.
26.5.4. Hot Standby Parameter Reference

Various parameters have been mentioned above in Section 26.5.2 and Section 26.5.3.

On the primary, parameters wal_level and vacuum_defer_cleanup_age can be used. max_standby_archive_delay and max_standby_streaming_delay have no effect if set on the primary.

On the standby, parameters hot_standby, max_standby_archive_delay and max_standby_streaming_delay can be used. vacuum_defer_cleanup_age has no effect as long as the server remains in standby mode, though it will become relevant if the standby becomes primary.

26.5.5. Caveats

There are several limitations of Hot Standby. These can and probably will be fixed in future releases:

- Full knowledge of running transactions is required before snapshots can be taken. Transactions that use large numbers of subtransactions (currently greater than 64) will delay the start of read only connections until the completion of the longest running write transaction. If this situation occurs, explanatory messages will be sent to the server log.

- Valid starting points for standby queries are generated at each checkpoint on the master. If the standby is shut down while the master is in a shutdown state, it might not be possible to re-enter Hot Standby until the primary is started up, so that it generates further starting points in the WAL logs. This situation isn't a problem in the most common situations where it might happen. Generally, if the primary is shut down and not available anymore, that's likely due to a serious failure that requires the standby being converted to operate as the new primary anyway. And in situations where the primary is being intentionally taken down, coordinating to make sure the standby becomes the new primary smoothly is also standard procedure.

- At the end of recovery, AccessExclusiveLocks held by prepared transactions will require twice the normal number of lock table entries. If you plan on running either a large number of concurrent prepared transactions that normally take AccessExclusiveLocks, or you plan on having one large transaction that takes many AccessExclusiveLocks, you are advised to select a larger value of max_locks_per_transaction, perhaps as much as twice the value of the parameter on the primary server. You need not consider this at all if your setting of max_prepared_transactions is 0.

- The Serializable transaction isolation level is not yet available in hot standby. (See Section 13.2.3 and Section 13.4.1 for details.) An attempt to set a transaction to the serializable isolation level in hot standby mode will generate an error.
Chapter 27. Recovery Configuration

This chapter describes the settings available in the `recovery.conf` file. They apply only for the duration of the recovery. They must be reset for any subsequent recovery you wish to perform. They cannot be changed once recovery has begun.

Settings in `recovery.conf` are specified in the format `name = 'value'`. One parameter is specified per line. Hash marks (`#`) designate the rest of the line as a comment. To embed a single quote in a parameter value, write two quotes (`"`).

A sample file, `share/recovery.conf.sample`, is provided in the installation's `share/` directory.

27.1. Archive Recovery Settings

`restore_command(string)`

The local shell command to execute to retrieve an archived segment of the WAL file series. This parameter is required for archive recovery, but optional for streaming replication. Any `%f` in the string is replaced by the name of the file to retrieve from the archive, and any `%p` is replaced by the copy destination path name on the server. (The path name is relative to the current working directory, i.e., the cluster's data directory.) Any `%r` is replaced by the name of the file containing the last valid restart point. That is the earliest file that must be kept to allow a restore to be restartable, so this information can be used to truncate the archive to just the minimum required to support restarting from the current restore. `%r` is typically only used by warm-standby configurations (see Section 26.2). Write `%%` to embed an actual `%` character.

It is important for the command to return a zero exit status only if it succeeds. The command will be asked for file names that are not present in the archive; it must return nonzero when so asked. Examples:

```
restore_command = 'cp /mnt/server/archivedir/%f "%p"'
restore_command = 'copy "C:\server\archivedir\%f" "%p"'  # Windows
```

An exception is that if the command was terminated by a signal (other than SIGTERM, which is used as part of a database server shutdown) or an error by the shell (such as command not found), then recovery will abort and the server will not start up.

`archive_cleanup_command(string)`

This optional parameter specifies a shell command that will be executed at every restartpoint. The purpose of `archive_cleanup_command` is to provide a mechanism for cleaning up old archived WAL files that are no longer needed by the standby server. Any `%r` is replaced by the name of the file containing the last valid restart point. That is the earliest file that must be kept to allow a restore to be restartable, and so all files earlier than `%r` may be safely removed. This information can be used to truncate the archive to just the minimum required to support restart from the current restore. The `pg_archivecleanup` module is often used in `archive_cleanup_command` for single-standby configurations, for example:

```
archive_cleanup_command = 'pg_archivecleanup /mnt/server/archivedir %r'
```

Note however that if multiple standby servers are restoring from the same archive directory, you will need to ensure that you do not delete WAL files until they are no longer needed by any of the servers. `archive_cleanup_command` would typically be used in a warm-standby configuration (see Section 26.2). Write `%%` to embed an actual `%` character in the command.
If the command returns a nonzero exit status then a warning log message will be written. An exception is that if the command was terminated by a signal or an error by the shell (such as command not found), a fatal error will be raised.

recovery_end_command(string)

This parameter specifies a shell command that will be executed once only at the end of recovery. This parameter is optional. The purpose of the recovery_end_command is to provide a mechanism for cleanup following replication or recovery. Any `%r` is replaced by the name of the file containing the last valid restart point, like in archive_cleanup_command.

If the command returns a nonzero exit status then a warning log message will be written and the database will proceed to start up anyway. An exception is that if the command was terminated by a signal or an error by the shell (such as command not found), the database will not proceed with startup.

## 27.2. Recovery Target Settings

By default, recovery will recover to the end of the WAL log. The following parameters can be used to specify an earlier stopping point. At most one of recovery_target, recovery_target_lsn, recovery_target_name, recovery_target_time, or recovery_target_xid can be used; if more than one of these is specified in the configuration file, the last entry will be used.

recovery_target = 'immediate'

This parameter specifies that recovery should end as soon as a consistent state is reached, i.e. as early as possible. When restoring from an online backup, this means the point where taking the backup ended.

Technically, this is a string parameter, but 'immediate' is currently the only allowed value.

recovery_target_name(string)

This parameter specifies the named restore point (created with pg_create_reSTORE_POINT()) to which recovery will proceed.

recovery_target_time(timestamp)

This parameter specifies the time stamp up to which recovery will proceed. The precise stopping point is also influenced by recovery_target_inclusive.

recovery_target_xid(string)

This parameter specifies the transaction ID up to which recovery will proceed. Keep in mind that while transaction IDs are assigned sequentially at transaction start, transactions can complete in a different numeric order. The transactions that will be recovered are those that committed before (and optionally including) the specified one. The precise stopping point is also influenced by recovery_target_inclusive.

recovery_target_lsn(pg_lsn)

This parameter specifies the LSN of the write-ahead log location up to which recovery will proceed. The precise stopping point is also influenced by recovery_target_inclusive. This parameter is parsed using the system data type pg_lsn.

The following options further specify the recovery target, and affect what happens when the target is reached:

recovery_target_inclusive(boolean)

Specifies whether to stop just after the specified recovery target (true), or just before the recovery target (false). Applies when recovery_target_lsn, recovery_target_time, or recovery_tar-
Recovery Configuration

get_xid is specified. This setting controls whether transactions having exactly the target WAL location (LSN), commit time, or transaction ID, respectively, will be included in the recovery. Default is true.

recovery_target_timeline (string)

Specifies recovering into a particular timeline. The default is to recover along the same timeline that was current when the base backup was taken. Setting this to latest recovers to the latest timeline found in the archive, which is useful in a standby server. Other than that you only need to set this parameter in complex re-recovery situations, where you need to return to a state that itself was reached after a point-in-time recovery. See Section 25.3.5 for discussion.

recovery_target_action (enum)

Specifies what action the server should take once the recovery target is reached. The default is pause, which means recovery will be paused. promote means the recovery process will finish and the server will start to accept connections. Finally shutdown will stop the server after reaching the recovery target.

The intended use of the pause setting is to allow queries to be executed against the database to check if this recovery target is the most desirable point for recovery. The paused state can be resumed by using pg_wal_replay_resume() (see Table 9.81), which then causes recovery to end. If this recovery target is not the desired stopping point, then shut down the server, change the recovery target settings to a later target and restart to continue recovery.

The shutdown setting is useful to have the instance ready at the exact replay point desired. The instance will still be able to replay more WAL records (and in fact will have to replay WAL records since the last checkpoint next time it is started).

Note that because recovery.conf will not be renamed when recovery_target_action is set to shutdown, any subsequent start will end with immediate shutdown unless the configuration is changed or the recovery.conf file is removed manually.

This setting has no effect if no recovery target is set. If hot_standby is not enabled, a setting of pause will act the same as shutdown.

27.3. Standby Server Settings

standby_mode (boolean)

Specifies whether to start the PostgreSQL server as a standby. If this parameter is on, the server will not stop recovery when the end of archived WAL is reached, but will keep trying to continue recovery by fetching new WAL segments using restore_command and/or by connecting to the primary server as specified by the primary_conninfo setting.

primary_conninfo (string)

Specifies a connection string to be used for the standby server to connect with the primary. This string is in the format described in Section 34.1.1. If any option is unspecified in this string, then the corresponding environment variable (see Section 34.14) is checked. If the environment variable is not set either, then defaults are used.

The connection string should specify the host name (or address) of the primary server, as well as the port number if it is not the same as the standby server's default. Also specify a user name corresponding to a suitably-privileged role on the primary (see Section 26.2.5.1). A password needs to be provided too, if the primary demands password authentication. It can be provided in the primary_conninfo string, or in a separate ~/.pgpass file on the standby server (use replication as the database name). Do not specify a database name in the primary_conninfo string.
Recovery Configuration

This setting has no effect if standby_mode is off.

primary_slot_name (string)

Optionally specifies an existing replication slot to be used when connecting to the primary via streaming replication to control resource removal on the upstream node (see Section 26.2.6). This setting has no effect if primary_conninfo is not set.

trigger_file (string)

Specifies a trigger file whose presence ends recovery in the standby. Even if this value is not set, you can still promote the standby using pg_ctl promote. This setting has no effect if standby_mode is off.

recovery_min_apply_delay (integer)

By default, a standby server restores WAL records from the primary as soon as possible. It may be useful to have a time-delayed copy of the data, offering opportunities to correct data loss errors. This parameter allows you to delay recovery by a fixed period of time, measured in milliseconds if no unit is specified. For example, if you set this parameter to 5min, the standby will replay each transaction commit only when the system time on the standby is at least five minutes past the commit time reported by the master.

It is possible that the replication delay between servers exceeds the value of this parameter, in which case no delay is added. Note that the delay is calculated between the WAL time stamp as written on master and the current time on the standby. Delays in transfer because of network lag or cascading replication configurations may reduce the actual wait time significantly. If the system clocks on master and standby are not synchronized, this may lead to recovery applying records earlier than expected; but that is not a major issue because useful settings of this parameter are much larger than typical time deviations between servers.

The delay occurs only on WAL records for transaction commits. Other records are replayed as quickly as possible, which is not a problem because MVCC visibility rules ensure their effects are not visible until the corresponding commit record is applied.

The delay occurs once the database in recovery has reached a consistent state, until the standby is promoted or triggered. After that the standby will end recovery without further waiting.

This parameter is intended for use with streaming replication deployments; however, if the parameter is specified it will be honored in all cases. hot_standby_feedback will be delayed by use of this feature which could lead to bloat on the master; use both together with care.

Warning

Synchronous replication is affected by this setting when synchronous_commit is set to remote_apply; every COMMIT will need to waits to be applied.
Chapter 28. Monitoring Database Activity

A database administrator frequently wonders, “What is the system doing right now?” This chapter discusses how to find that out.

Several tools are available for monitoring database activity and analyzing performance. Most of this chapter is devoted to describing PostgreSQL’s statistics collector, but one should not neglect regular Unix monitoring programs such as `ps`, `top`, `iostat`, and `vmstat`. Also, once one has identified a poorly-performing query, further investigation might be needed using PostgreSQL’s `EXPLAIN` command. Section 14.1 discusses `EXPLAIN` and other methods for understanding the behavior of an individual query.

### 28.1. Standard Unix Tools

On most Unix platforms, PostgreSQL modifies its command title as reported by `ps`, so that individual server processes can readily be identified. A sample display is

```
$ ps auxww | grep ^postgres
postgres  15551  0.0  0.1  57536  7132 pts/0    S    18:02   0:00
postgres -i
postgres  15554  0.0  0.0  57536  1184 ?        Ss   18:02   0:00
postgres: background writer
postgres  15555  0.0  0.0  57536   916 ?        Ss   18:02   0:00
postgres: checkpointer
postgres  15556  0.0  0.0  57536   916 ?        Ss   18:02   0:00
postgres: walwriter
postgres  15557  0.0  0.0  58504  2244 ?        Ss   18:02   0:00
postgres: autovacuum launcher
postgres  15558  0.0  0.0  17512  1068 ?        Ss   18:02   0:00
postgres: stats collector
postgres  15582  0.0  0.0  58772  3080 ?        Ss   18:04   0:00
postgres: joe runbug 127.0.0.1 idle
postgres  15606  0.0  0.0  58772  3052 ?        Ss   18:07   0:00
postgres: tgl regression [local] SELECT waiting
postgres  15610  0.0  0.0  58772  3056 ?        Ss   18:07   0:00
postgres: tgl regression [local] idle in transaction
```

(The appropriate invocation of `ps` varies across different platforms, as do the details of what is shown. This example is from a recent Linux system.) The first process listed here is the master server process. The command arguments shown for it are the same ones used when it was launched. The next five processes are background worker processes automatically launched by the master process. (The “stats collector” process will not be present if you have set the system not to start the statistics collector; likewise the “autovacuum launcher” process can be disabled.) Each of the remaining processes is a server process handling one client connection. Each such process sets its command line display in the form

```
postgres: user database host activity
```

The user, database, and (client) host items remain the same for the life of the client connection, but the activity indicator changes. The activity can be `idle` (i.e., waiting for a client command), `idle in transaction` (waiting for client inside a `BEGIN` block), or a command type name such as `SELECT`. Also, `waiting` is appended if the server process is presently waiting on a lock held by another session. In the above example we can infer that process 15606 is waiting for process 15610 to
Monitoring Database Activity

complete its transaction and thereby release some lock. (Process 15610 must be the blocker, because there is no other active session. In more complicated cases it would be necessary to look into the pg_locks system view to determine who is blocking whom.)

If cluster_name has been configured the cluster name will also be shown in ps output:

$ psql -c 'SHOW cluster_name'
    cluster_name
    ---------
    server1
    (1 row)

$ ps aux|grep server1
    postgres  27093  0.0  0.0  30096  2752 ?        Ss   11:34   0:00
    postgres: server1: background writer
    ...

If you have turned off update_process_title then the activity indicator is not updated; the process title is set only once when a new process is launched. On some platforms this saves a measurable amount of per-command overhead; on others it’s insignificant.

Tip

Solaris requires special handling. You must use /usr/ucb/ps, rather than /bin/ps. You also must use two w flags, not just one. In addition, your original invocation of the postgres command must have a shorter ps status display than that provided by each server process. If you fail to do all three things, the ps output for each server process will be the original postgres command line.

28.2. The Statistics Collector

PostgreSQL’s statistics collector is a subsystem that supports collection and reporting of information about server activity. Presently, the collector can count accesses to tables and indexes in both disk-block and individual-row terms. It also tracks the total number of rows in each table, and information about vacuum and analyze actions for each table. It can also count calls to user-defined functions and the total time spent in each one.

PostgreSQL also supports reporting dynamic information about exactly what is going on in the system right now, such as the exact command currently being executed by other server processes, and which other connections exist in the system. This facility is independent of the collector process.

28.2.1. Statistics Collection Configuration

Since collection of statistics adds some overhead to query execution, the system can be configured to collect or not collect information. This is controlled by configuration parameters that are normally set in postgresql.conf. (See Chapter 19 for details about setting configuration parameters.)

The parameter track_activities enables monitoring of the current command being executed by any server process.

The parameter track_counts controls whether statistics are collected about table and index accesses.

The parameter track_functions enables tracking of usage of user-defined functions.

The parameter track_io_timing enables monitoring of block read and write times.

Normally these parameters are set in postgresql.conf so that they apply to all server processes, but it is possible to turn them on or off in individual sessions using the SET command. (To prevent
 ordinary users from hiding their activity from the administrator, only superusers are allowed to change these parameters with \texttt{SET}.)

The statistics collector transmits the collected information to other PostgreSQL processes through temporary files. These files are stored in the directory named by the \texttt{stats_temp_directory} parameter, \texttt{pg_stat_tmp} by default. For better performance, \texttt{stats_temp_directory} can be pointed at a RAM-based file system, decreasing physical I/O requirements. When the server shuts down cleanly, a permanent copy of the statistics data is stored in the \texttt{pg_stat} subdirectory, so that statistics can be retained across server restarts. When recovery is performed at server start (e.g. after immediate shutdown, server crash, and point-in-time recovery), all statistics counters are reset.

### 28.2.2. Viewing Statistics

Several predefined views, listed in Table 28.1, are available to show the current state of the system. There are also several other views, listed in Table 28.2, available to show the results of statistics collection. Alternatively, one can build custom views using the underlying statistics functions, as discussed in Section 28.2.3.

When using the statistics to monitor collected data, it is important to realize that the information does not update instantaneously. Each individual server process transmits new statistical counts to the collector just before going idle; so a query or transaction still in progress does not affect the displayed totals. Also, the collector itself emits a new report at most once per \texttt{PGSTAT_STAT_INTERVAL} milliseconds (500 ms unless altered while building the server). So the displayed information lags behind actual activity. However, current-query information collected by \texttt{track_activities} is always up-to-date.

Another important point is that when a server process is asked to display any of these statistics, it first fetches the most recent report emitted by the collector process and then continues to use this snapshot for all statistical views and functions until the end of its current transaction. So the statistics will show static information as long as you continue the current transaction. Similarly, information about the current queries of all sessions is collected when any such information is first requested within a transaction, and the same information will be displayed throughout the transaction. This is a feature, not a bug, because it allows you to perform several queries on the statistics and correlate the results without worrying that the numbers are changing underneath you. But if you want to see new results with each query, be sure to do the queries outside any transaction block. Alternatively, you can invoke \texttt{pg_stat_clear_snapshot()}, which will discard the current transaction's statistics snapshot (if any). The next use of statistical information will cause a new snapshot to be fetched.

A transaction can also see its own statistics (as yet untransmitted to the collector) in the views \texttt{pg_stat_xact_all_tables}, \texttt{pg_stat_xact_sys_tables}, \texttt{pg_stat_xact_user_tables}, and \texttt{pg_stat_xact_user_functions}. These numbers do not act as stated above; instead they update continuously throughout the transaction.

#### Table 28.1. Dynamic Statistics Views

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{pg_stat_activity}</td>
<td>One row per server process, showing information related to the current activity of that process, such as state and current query. See \texttt{pg_stat_activity} for details.</td>
</tr>
<tr>
<td>\texttt{pg_stat_replication}</td>
<td>One row per WAL sender process, showing statistics about replication to that sender's connected standby server. See \texttt{pg_stat_replication} for details.</td>
</tr>
<tr>
<td>\texttt{pg_stat_wal_receiver}</td>
<td>Only one row, showing statistics about the WAL receiver from that receiver's connected server. See \texttt{pg_stat_wal_receiver} for details.</td>
</tr>
</tbody>
</table>
### Monitoring Database Activity

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_stat_subscription</code></td>
<td>At least one row per subscription, showing information about the subscription workers. See <code>pg_stat_subscription</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_ssl</code></td>
<td>One row per connection (regular and replication), showing information about SSL used on this connection. See <code>pg_stat_ssl</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_progress_vacuum</code></td>
<td>One row for each backend (including autovacuum worker processes) running VACUUM, showing current progress. See Section 28.4.1.</td>
</tr>
</tbody>
</table>

#### Table 28.2. Collected Statistics Views

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_stat_archiver</code></td>
<td>One row only, showing statistics about the WAL archiver process’s activity. See <code>pg_stat_archiver</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_bgwriter</code></td>
<td>One row only, showing statistics about the background writer process’s activity. See <code>pg_stat_bgwriter</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_database</code></td>
<td>One row per database, showing database-wide statistics. See <code>pg_stat_database</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_database_conflicts</code></td>
<td>One row per database, showing database-wide statistics about query cancels due to conflict with recovery on standby servers. See <code>pg_stat_database_conflicts</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_all_tables</code></td>
<td>One row for each table in the current database, showing statistics about accesses to that specific table. See <code>pg_stat_all_tables</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_sys_tables</code></td>
<td>Same as <code>pg_stat_all_tables</code>, except that only system tables are shown.</td>
</tr>
<tr>
<td><code>pg_stat_user_tables</code></td>
<td>Same as <code>pg_stat_all_tables</code>, except that only user tables are shown.</td>
</tr>
<tr>
<td><code>pg_stat_xact_all_tables</code></td>
<td>Similar to <code>pg_stat_all_tables</code>, but counts actions taken so far within the current transaction (which are not yet included in <code>pg_stat_all_tables</code> and related views). The columns for numbers of live and dead rows and vacuum and analyze actions are not present in this view.</td>
</tr>
<tr>
<td><code>pg_stat_xact_sys_tables</code></td>
<td>Same as <code>pg_stat_xact_all_tables</code>, except that only system tables are shown.</td>
</tr>
<tr>
<td><code>pg_stat_xact_user_tables</code></td>
<td>Same as <code>pg_stat_xact_all_tables</code>, except that only user tables are shown.</td>
</tr>
<tr>
<td><code>pg_stat_all_indexes</code></td>
<td>One row for each index in the current database, showing statistics about accesses to that specific index. See <code>pg_stat_all_indexes</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_sys_indexes</code></td>
<td>Same as <code>pg_stat_all_indexes</code>, except that only indexes on system tables are shown.</td>
</tr>
<tr>
<td><code>pg_stat_user_indexes</code></td>
<td>Same as <code>pg_stat_all_indexes</code>, except that only indexes on user tables are shown.</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_statio_all_tables</td>
<td>One row for each table in the current database, showing statistics about I/O on that specific table. See pg_statio_all_tables for details.</td>
</tr>
<tr>
<td>pg_statio_sys_tables</td>
<td>Same as pg_statio_all_tables, except that only system tables are shown.</td>
</tr>
<tr>
<td>pg_statio_user_tables</td>
<td>Same as pg_statio_all_tables, except that only user tables are shown.</td>
</tr>
<tr>
<td>pg_statio_all_indexes</td>
<td>One row for each index in the current database, showing statistics about I/O on that specific index. See pg_statio_all_indexes for details.</td>
</tr>
<tr>
<td>pg_statio_sys_indexes</td>
<td>Same as pg_statio_all_indexes, except that only indexes on system tables are shown.</td>
</tr>
<tr>
<td>pg_statio_user_indexes</td>
<td>Same as pg_statio_all_indexes, except that only indexes on user tables are shown.</td>
</tr>
<tr>
<td>pg_statio_all_sequences</td>
<td>One row for each sequence in the current database, showing statistics about I/O on that specific sequence. See pg_statio_all_sequences for details.</td>
</tr>
<tr>
<td>pg_statio_sys_sequences</td>
<td>Same as pg_statio_all_sequences, except that only system sequences are shown. (Presently, no system sequences are defined, so this view is always empty.)</td>
</tr>
<tr>
<td>pg_statio_user_sequences</td>
<td>Same as pg_statio_all_sequences, except that only user sequences are shown.</td>
</tr>
<tr>
<td>pg_stat_user_functions</td>
<td>One row for each tracked function, showing statistics about executions of that function. See pg_stat_user_functions for details.</td>
</tr>
<tr>
<td>pg_stat_xact_user_functions</td>
<td>Similar to pg_stat_user_functions, but counts only calls during the current transaction (which are not yet included in pg_stat_user_functions).</td>
</tr>
</tbody>
</table>

The per-index statistics are particularly useful to determine which indexes are being used and how effective they are.

The pg_statio_ views are primarily useful to determine the effectiveness of the buffer cache. When the number of actual disk reads is much smaller than the number of buffer hits, then the cache is satisfying most read requests without invoking a kernel call. However, these statistics do not give the entire story: due to the way in which PostgreSQL handles disk I/O, data that is not in the PostgreSQL buffer cache might still reside in the kernel's I/O cache, and might therefore still be fetched without requiring a physical read. Users interested in obtaining more detailed information on PostgreSQL I/O behavior are advised to use the PostgreSQL statistics collector in combination with operating system utilities that allow insight into the kernel's handling of I/O.

**Table 28.3. pg_stat_activity View**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of the database this backend is connected to</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of the database this backend is connected to</td>
</tr>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of this backend</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>OID of the user logged into this backend</td>
</tr>
<tr>
<td>username</td>
<td>name</td>
<td>Name of the user logged into this backend</td>
</tr>
<tr>
<td>application_name</td>
<td>text</td>
<td>Name of the application that is connected to this backend</td>
</tr>
<tr>
<td>client_addr</td>
<td>inet</td>
<td>IP address of the client connected to this backend. If this field is null, it indicates either that the client is connected via a Unix socket on the server machine or that this is an internal process such as autovacuum.</td>
</tr>
<tr>
<td>client_hostname</td>
<td>text</td>
<td>Host name of the connected client, as reported by a reverse DNS lookup of client_addr. This field will only be non-null for IP connections, and only when log_hostname is enabled.</td>
</tr>
<tr>
<td>client_port</td>
<td>integer</td>
<td>TCP port number that the client is using for communication with this backend, or -1 if a Unix socket is used</td>
</tr>
<tr>
<td>backend_start</td>
<td>timestamp with time zone</td>
<td>Time when this process was started. For client backends, this is the time the client connected to the server.</td>
</tr>
<tr>
<td>xact_start</td>
<td>timestamp with time zone</td>
<td>Time when this process’ current transaction was started, or null if no transaction is active. If the current query is the first of its transaction, this column is equal to the query_start column.</td>
</tr>
<tr>
<td>query_start</td>
<td>timestamp with time zone</td>
<td>Time when the currently active query was started, or if state is not active, when the last query was started</td>
</tr>
<tr>
<td>state_change</td>
<td>timestamp with time zone</td>
<td>Time when the state was last changed</td>
</tr>
<tr>
<td>wait_event_type</td>
<td>text</td>
<td>The type of event for which the backend is waiting, if any; otherwise NULL. Possible values are:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• LWLock: The backend is waiting for a lightweight lock. Each such lock protects a particular data structure in shared memory. wait_event will contain a name identifying the purpose of the lightweight lock. (Some locks have specific names; others are part of</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>a group of locks each with a similar purpose.)</td>
</tr>
<tr>
<td></td>
<td>Lock</td>
<td>The backend is waiting for a heavyweight lock. Heavyweight locks, also known as lock manager locks or simply locks, primarily protect SQL-visible objects such as tables. However, they are also used to ensure mutual exclusion for certain internal operations such as relation extension. wait_event will identify the type of lock awaited.</td>
</tr>
<tr>
<td></td>
<td>BufferPin</td>
<td>The server process is waiting to access to a data buffer during a period when no other process can be examining that buffer. Buffer pin waits can be protracted if another process holds an open cursor which last read data from the buffer in question.</td>
</tr>
<tr>
<td></td>
<td>Activity</td>
<td>The server process is idle. This is used by system processes waiting for activity in their main processing loop. wait_event will identify the specific wait point.</td>
</tr>
<tr>
<td></td>
<td>Extension</td>
<td>The server process is waiting for activity in an extension module. This category is useful for modules to track custom waiting points.</td>
</tr>
<tr>
<td></td>
<td>Client</td>
<td>The server process is waiting for some activity on a socket from user applications, and that the server expects something to happen that is independent from its internal processes. wait_event will identify the specific wait point.</td>
</tr>
<tr>
<td></td>
<td>IPC</td>
<td>The server process is waiting for some activity from another process in the server. wait_event will identify the specific wait point.</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------</td>
<td>------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>Timeout</strong>: The server process is waiting for a timeout to expire. <em>wait_event</em> will identify the specific wait point.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>IO</strong>: The server process is waiting for an IO to complete. <em>wait_event</em> will identify the specific wait point.</td>
</tr>
<tr>
<td>wait_event</td>
<td>text</td>
<td>Wait event name if backend is currently waiting, otherwise NULL. See Table 28.4 for details.</td>
</tr>
<tr>
<td>state</td>
<td>text</td>
<td>Current overall state of this backend. Possible values are:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>active</strong>: The backend is executing a query.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>idle</strong>: The backend is waiting for a new client command.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>idle in transaction</strong>: The backend is in a transaction, but is not currently executing a query.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>idle in transaction (aborted)</strong>: This state is similar to <strong>idle in transaction</strong>, except one of the statements in the transaction caused an error.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>fastpath function call</strong>: The backend is executing a fast-path function.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• <strong>disabled</strong>: This state is reported if track_activities is disabled in this backend.</td>
</tr>
<tr>
<td>backend_xid</td>
<td>xid</td>
<td>Top-level transaction identifier of this backend, if any.</td>
</tr>
<tr>
<td>backend_xmin</td>
<td>xid</td>
<td>The current backend's xmin horizon.</td>
</tr>
<tr>
<td>query</td>
<td>text</td>
<td>Text of this backend's most recent query. If <strong>state</strong> is <strong>active</strong> this field shows the currently executing query. In all other states, it shows the last query that was executed. By default the query text is truncated at 1024 characters; this value can be changed via the parameter <strong>track_activity_query_size</strong>.</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>backend_type</td>
<td>text</td>
<td>Type of current backend. Possible types are autovacuum launcher, autovacuum worker, background worker, background writer, client backend, checkpoint, startup, walreceiver, walsender and walwriter.</td>
</tr>
</tbody>
</table>

The `pg_stat_activity` view will have one row per server process, showing information related to the current activity of that process.

**Note**

The `wait_event` and `state` columns are independent. If a backend is in the active state, it may or may not be waiting on some event. If the state is active and `wait_event` is non-null, it means that a query is being executed, but is being blocked somewhere in the system.

**Table 28.4. `wait_event` Description**

<table>
<thead>
<tr>
<th>Wait Event Type</th>
<th>Wait Event Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>LWLock</td>
<td>ShmemIndexLock</td>
<td>Waiting to find or allocate space in shared memory.</td>
</tr>
<tr>
<td></td>
<td>OidGenLock</td>
<td>Waiting to allocate or assign an OID.</td>
</tr>
<tr>
<td></td>
<td>XidGenLock</td>
<td>Waiting to allocate or assign a transaction id.</td>
</tr>
<tr>
<td></td>
<td>ProcArrayLock</td>
<td>Waiting to get a snapshot or clearing a transaction id at transaction end.</td>
</tr>
<tr>
<td></td>
<td>SINvalReadLock</td>
<td>Waiting to retrieve or remove messages from shared invalidation queue.</td>
</tr>
<tr>
<td></td>
<td>SINvalWriteLock</td>
<td>Waiting to add a message in shared invalidation queue.</td>
</tr>
<tr>
<td></td>
<td>WALBufMappingLock</td>
<td>Waiting to replace a page in WAL buffers.</td>
</tr>
<tr>
<td></td>
<td>WALWriteLock</td>
<td>Waiting for WAL buffers to be written to disk.</td>
</tr>
<tr>
<td></td>
<td>ControlFileLock</td>
<td>Waiting to read or update the control file or creation of a new WAL file.</td>
</tr>
<tr>
<td></td>
<td>CheckpointLock</td>
<td>Waiting to perform checkpoint.</td>
</tr>
<tr>
<td></td>
<td>CLogControlLock</td>
<td>Waiting to read or update transaction status.</td>
</tr>
<tr>
<td></td>
<td>SubtransControlLock</td>
<td>Waiting to read or update subtransaction information.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------------------</td>
<td>----------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>MultiXactGenLock</td>
<td>Waiting to read or update shared multixact state.</td>
<td></td>
</tr>
<tr>
<td>MultiXactOffsetControlLock</td>
<td>Waiting to read or update multixact offset mappings.</td>
<td></td>
</tr>
<tr>
<td>MultiXactMemberControlLock</td>
<td>Waiting to read or update multixact member mappings.</td>
<td></td>
</tr>
<tr>
<td>RelCacheInitLock</td>
<td>Waiting to read or write relation cache initialization file.</td>
<td></td>
</tr>
<tr>
<td>CheckpointerCommLock</td>
<td>Waiting to manage fsync requests.</td>
<td></td>
</tr>
<tr>
<td>TwoPhaseStateLock</td>
<td>Waiting to read or update the state of prepared transactions.</td>
<td></td>
</tr>
<tr>
<td>TablespaceCreateLock</td>
<td>Waiting to create or drop the tablespace.</td>
<td></td>
</tr>
<tr>
<td>BtreeVacuumLock</td>
<td>Waiting to read or update vacuum-related information for a B-tree index.</td>
<td></td>
</tr>
<tr>
<td>AddinShmemInitLock</td>
<td>Waiting to manage space allocation in shared memory.</td>
<td></td>
</tr>
<tr>
<td>AutovacuumLock</td>
<td>Autovacuum worker or launcher waiting to update or read the current state of autovacuum workers.</td>
<td></td>
</tr>
<tr>
<td>AutovacuumScheduleLock</td>
<td>Waiting to ensure that the table it has selected for a vacuum still needs vacuuming.</td>
<td></td>
</tr>
<tr>
<td>SyncScanLock</td>
<td>Waiting to get the start location of a scan on a table for synchronized scans.</td>
<td></td>
</tr>
<tr>
<td>RelationMappingLock</td>
<td>Waiting to update the relation map file used to store catalog to filenode mapping.</td>
<td></td>
</tr>
<tr>
<td>AsyncCtlLock</td>
<td>Waiting to read or update shared notification state.</td>
<td></td>
</tr>
<tr>
<td>AsyncQueueLock</td>
<td>Waiting to read or update notification messages.</td>
<td></td>
</tr>
<tr>
<td>SerializableXactHashLock</td>
<td>Waiting to retrieve or store information about serializable transactions.</td>
<td></td>
</tr>
<tr>
<td>SerializableFinishedListLock</td>
<td>Waiting to access the list of finished serializable transactions.</td>
<td></td>
</tr>
<tr>
<td>SerializablePredicateLockListLock</td>
<td>Waiting to perform an operation on a list of locks held by serializable transactions.</td>
<td></td>
</tr>
<tr>
<td>OldSerXidLock</td>
<td>Waiting to read or record conflicting serializable transactions.</td>
<td></td>
</tr>
<tr>
<td>SyncRepLock</td>
<td>Waiting to read or update information about synchronous replicas.</td>
<td></td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------------</td>
<td>----------------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>BackgroundWorkerLock</td>
<td>Waiting to read or update background worker state.</td>
<td></td>
</tr>
<tr>
<td>DynamicSharedMemoryControlLock</td>
<td>Waiting to read or update dynamic shared memory state.</td>
<td></td>
</tr>
<tr>
<td>AutoFileLock</td>
<td>Waiting to update the postgresql.auto.conf file.</td>
<td></td>
</tr>
<tr>
<td>ReplicationSlotAllocationLock</td>
<td>Waiting to allocate or free a replication slot.</td>
<td></td>
</tr>
<tr>
<td>ReplicationSlotControlLock</td>
<td>Waiting to read or update replication slot state.</td>
<td></td>
</tr>
<tr>
<td>CommitTsControlLock</td>
<td>Waiting to read or update transaction commit timestamps.</td>
<td></td>
</tr>
<tr>
<td>CommitTsLock</td>
<td>Waiting to read or update the last value set for the transaction timestamp.</td>
<td></td>
</tr>
<tr>
<td>ReplicationOriginLock</td>
<td>Waiting to setup, drop or use replication origin.</td>
<td></td>
</tr>
<tr>
<td>MultiXactTruncationLock</td>
<td>Waiting to read or truncate multixact information.</td>
<td></td>
</tr>
<tr>
<td>OldSnapshot-TimeMapLock</td>
<td>Waiting to read or update old snapshot control information.</td>
<td></td>
</tr>
<tr>
<td>BackendRandomLock</td>
<td>Waiting to generate a random number.</td>
<td></td>
</tr>
<tr>
<td>LogicalRepWorkerLock</td>
<td>Waiting for action on logical replication worker to finish.</td>
<td></td>
</tr>
<tr>
<td>CLogTruncationLock</td>
<td>Waiting to truncate the write-ahead log or waiting for write-ahead log truncation to finish.</td>
<td></td>
</tr>
<tr>
<td>clog</td>
<td>Waiting for I/O on a clog (transaction status) buffer.</td>
<td></td>
</tr>
<tr>
<td>commit_timestamp</td>
<td>Waiting for I/O on commit timestamp buffer.</td>
<td></td>
</tr>
<tr>
<td>subtrans</td>
<td>Waiting for I/O a subtransaction buffer.</td>
<td></td>
</tr>
<tr>
<td>multixact_offset</td>
<td>Waiting for I/O on a multixact offset buffer.</td>
<td></td>
</tr>
<tr>
<td>multixact_member</td>
<td>Waiting for I/O on a multixact_member buffer.</td>
<td></td>
</tr>
<tr>
<td>async</td>
<td>Waiting for I/O on an async (notify) buffer.</td>
<td></td>
</tr>
<tr>
<td>oldserxid</td>
<td>Waiting for I/O on an oldserxid buffer.</td>
<td></td>
</tr>
<tr>
<td>wal_insert</td>
<td>Waiting to insert WAL into a memory buffer.</td>
<td></td>
</tr>
<tr>
<td>buffer_content</td>
<td>Waiting to read or write a data page in memory.</td>
<td></td>
</tr>
<tr>
<td>buffer_io</td>
<td>Waiting for I/O on a data page.</td>
<td></td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>replication_origin</td>
<td>Waiting to read or update the replication progress.</td>
<td></td>
</tr>
<tr>
<td>replication_slot_io</td>
<td>Waiting for I/O on a replication slot.</td>
<td></td>
</tr>
<tr>
<td>proc</td>
<td>Waiting to read or update the fast-path lock information.</td>
<td></td>
</tr>
<tr>
<td>buffer_mapping</td>
<td>Waiting to associate a data block with a buffer in the buffer pool.</td>
<td></td>
</tr>
<tr>
<td>lock_manager</td>
<td>Waiting to add or examine locks for backends, or waiting to join or exit a locking group (used by parallel query).</td>
<td></td>
</tr>
<tr>
<td>predicate_lock_manager</td>
<td>Waiting to add or examine predicate lock information.</td>
<td></td>
</tr>
<tr>
<td>parallel_query_dsa</td>
<td>Waiting for parallel query dynamic shared memory allocation lock.</td>
<td></td>
</tr>
<tr>
<td>tbm</td>
<td>Waiting for TBM shared iterator lock.</td>
<td></td>
</tr>
<tr>
<td>parallel_append</td>
<td>Waiting to choose the next subplan during Parallel Append plan execution.</td>
<td></td>
</tr>
<tr>
<td>parallel_hash_join</td>
<td>Waiting to allocate or exchange a chunk of memory or update counters during Parallel Hash plan execution.</td>
<td></td>
</tr>
<tr>
<td>relation</td>
<td>Waiting to acquire a lock on a relation.</td>
<td></td>
</tr>
<tr>
<td>extend</td>
<td>Waiting to extend a relation.</td>
<td></td>
</tr>
<tr>
<td>page</td>
<td>Waiting to acquire a lock on page of a relation.</td>
<td></td>
</tr>
<tr>
<td>tuple</td>
<td>Waiting to acquire a lock on a tuple.</td>
<td></td>
</tr>
<tr>
<td>transactionid</td>
<td>Waiting for a transaction to finish.</td>
<td></td>
</tr>
<tr>
<td>virtualxid</td>
<td>Waiting to acquire a virtual xid lock.</td>
<td></td>
</tr>
<tr>
<td>speculative token</td>
<td>Waiting to acquire a speculative insertion lock.</td>
<td></td>
</tr>
<tr>
<td>object</td>
<td>Waiting to acquire a lock on a non-relation database object.</td>
<td></td>
</tr>
<tr>
<td>userlock</td>
<td>Waiting to acquire a user lock.</td>
<td></td>
</tr>
<tr>
<td>advisory</td>
<td>Waiting to acquire an advisory user lock.</td>
<td></td>
</tr>
<tr>
<td>BufferPin</td>
<td>BufferPin</td>
<td>Waiting to acquire a pin on a buffer.</td>
</tr>
<tr>
<td>Activity</td>
<td>ArchiverMain</td>
<td>Waiting in main loop of the archiver process.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>---------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>AutoVacuumMain</td>
<td>Waiting in main loop of autovacuum launcher process.</td>
<td></td>
</tr>
<tr>
<td>BgWriterHibernate</td>
<td>Waiting in background writer process, hibernating.</td>
<td></td>
</tr>
<tr>
<td>BgWriterMain</td>
<td>Waiting in main loop of background writer process background worker.</td>
<td></td>
</tr>
<tr>
<td>CheckpointerMain</td>
<td>Waiting in main loop of checkpoint process.</td>
<td></td>
</tr>
<tr>
<td>LogicalApplyMain</td>
<td>Waiting in main loop of logical apply process.</td>
<td></td>
</tr>
<tr>
<td>LogicalLauncherMain</td>
<td>Waiting in main loop of logical launcher process.</td>
<td></td>
</tr>
<tr>
<td>PgStatMain</td>
<td>Waiting in main loop of the statistics collector process.</td>
<td></td>
</tr>
<tr>
<td>RecoveryWalAll</td>
<td>Waiting for WAL from any kind of source (local, archive or stream) at recovery.</td>
<td></td>
</tr>
<tr>
<td>RecoveryWalStream</td>
<td>Waiting for WAL from a stream at recovery.</td>
<td></td>
</tr>
<tr>
<td>SysLoggerMain</td>
<td>Waiting in main loop of syslogger process.</td>
<td></td>
</tr>
<tr>
<td>WalReceiverMain</td>
<td>Waiting in main loop of WAL receiver process.</td>
<td></td>
</tr>
<tr>
<td>WalSenderMain</td>
<td>Waiting in main loop of WAL sender process.</td>
<td></td>
</tr>
<tr>
<td>WalWriterMain</td>
<td>Waiting in main loop of WAL writer process.</td>
<td></td>
</tr>
<tr>
<td>Client</td>
<td>ClientRead</td>
<td>Waiting to read data from the client.</td>
</tr>
<tr>
<td></td>
<td>ClientWrite</td>
<td>Waiting to write data to the client.</td>
</tr>
<tr>
<td></td>
<td>LibPQWalReceiverConnect</td>
<td>Waiting in WAL receiver to establish connection to remote server.</td>
</tr>
<tr>
<td></td>
<td>LibPQWalReceiverReceive</td>
<td>Waiting in WAL receiver to receive data from remote server.</td>
</tr>
<tr>
<td></td>
<td>SSLOpenServer</td>
<td>Waiting for SSL while attempting connection.</td>
</tr>
<tr>
<td></td>
<td>WalReceiverWaitStart</td>
<td>Waiting for startup process to send initial data for streaming replication.</td>
</tr>
<tr>
<td></td>
<td>WalSenderWaitForWAL</td>
<td>Waiting for WAL to be flushed in WAL sender process.</td>
</tr>
<tr>
<td></td>
<td>WalSenderWriteData</td>
<td>Waiting for any activity when processing replies from WAL receiver in WAL sender process.</td>
</tr>
<tr>
<td>Extension</td>
<td>Extension</td>
<td>Waiting in an extension.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>---------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>IPC</td>
<td>BgWorkerShutdown</td>
<td>Waiting for background worker to shut down.</td>
</tr>
<tr>
<td></td>
<td>BgWorkerStartup</td>
<td>Waiting for background worker to start up.</td>
</tr>
<tr>
<td></td>
<td>BtreePage</td>
<td>Waiting for the page number needed to continue a parallel B-tree scan to become available.</td>
</tr>
<tr>
<td></td>
<td>ClogGroupUpdate</td>
<td>Waiting for group leader to update transaction status at transaction end.</td>
</tr>
<tr>
<td></td>
<td>ExecuteGather</td>
<td>Waiting for activity from child process when executing Gather node.</td>
</tr>
<tr>
<td></td>
<td>Hash/Batch/Allocating</td>
<td>Waiting for an elected Parallel Hash participant to allocate a hash table.</td>
</tr>
<tr>
<td></td>
<td>Hash/Batch/Electing</td>
<td>Electing a Parallel Hash participant to allocate a hash table.</td>
</tr>
<tr>
<td></td>
<td>Hash/Batch/Loading</td>
<td>Waiting for other Parallel Hash participants to finish loading a hash table.</td>
</tr>
<tr>
<td></td>
<td>Hash/Build/Allocating</td>
<td>Waiting for an elected Parallel Hash participant to allocate the initial hash table.</td>
</tr>
<tr>
<td></td>
<td>Hash/Build/Electing</td>
<td>Electing a Parallel Hash participant to allocate the initial hash table.</td>
</tr>
<tr>
<td></td>
<td>Hash/Build/HashingInner</td>
<td>Waiting for other Parallel Hash participants to finish hashing the inner relation.</td>
</tr>
<tr>
<td></td>
<td>Hash/Build/HashingOuter</td>
<td>Waiting for other Parallel Hash participants to finish partitioning the outer relation.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBatches/Allocating</td>
<td>Waiting for an elected Parallel Hash participant to allocate more batches.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBatches/Deciding</td>
<td>Electing a Parallel Hash participant to decide on future batch growth.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBatches/Electing</td>
<td>Electing a Parallel Hash participant to allocate more batches.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBatches/Finishing</td>
<td>Waiting for an elected Parallel Hash participant to decide on future batch growth.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBatches/Repartitioning</td>
<td>Waiting for other Parallel Hash participants to finishing repartitioning.</td>
</tr>
<tr>
<td></td>
<td>Hash/GrowBuckets/Allocating</td>
<td>Waiting for an elected Parallel Hash participant to finish allocating more buckets.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------------</td>
<td>-----------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>Hash/GrowBuckets/Electing</td>
<td>Electing a Parallel Hash particip-</td>
<td>waiting to elect a new participant to allocate more buckets.</td>
</tr>
<tr>
<td>Hash/GrowBuckets/Reinserting</td>
<td>Waiting for other Parallel Hash</td>
<td>waiting for other Parallel Hash participants to finish inserting tuples into new buckets.</td>
</tr>
<tr>
<td>LogicalSyncData</td>
<td>Waiting for logical replication</td>
<td>waiting for logical replication remote server to send data for initial table synchronization.</td>
</tr>
<tr>
<td>LogicalSyncStateChanged</td>
<td>Waiting for logical replication</td>
<td>waiting for logical replication remote server to change state.</td>
</tr>
<tr>
<td>MessageQueueInternal</td>
<td>Waiting for other process to be</td>
<td>waiting for other process to be attached in shared message queue.</td>
</tr>
<tr>
<td>MessageQueuePutMessage</td>
<td>Waiting to write a protocol mes-</td>
<td>waiting to write a protocol message to a shared message queue.</td>
</tr>
<tr>
<td>MessageQueueReceive</td>
<td>Waiting to receive bytes from a</td>
<td>waiting to receive bytes from a shared message queue.</td>
</tr>
<tr>
<td>MessageQueueSend</td>
<td>Waiting to send bytes to a shared</td>
<td>waiting to send bytes to a shared message queue.</td>
</tr>
<tr>
<td>ParallelBitmapScan</td>
<td>Waiting for parallel bitmap scan</td>
<td>waiting for parallel bitmap scan to become initialized.</td>
</tr>
<tr>
<td>ParallelCreateIndexScan</td>
<td>Waiting for parallel CREATE INDEX workers to finish heap scan.</td>
<td></td>
</tr>
<tr>
<td>ParallelFinish</td>
<td>Waiting for parallel workers to</td>
<td>waiting for parallel workers to finish computing.</td>
</tr>
<tr>
<td>ProcArrayGroupUpdate</td>
<td>Waiting for group leader to clear</td>
<td>waiting for group leader to clear transaction id at transaction end.</td>
</tr>
<tr>
<td>ReplicationOriginDrop</td>
<td>Waiting for a replication origin to become inactive to be dropped.</td>
<td></td>
</tr>
<tr>
<td>ReplicationSlotDrop</td>
<td>Waiting for a replication slot to become inactive to be dropped.</td>
<td></td>
</tr>
<tr>
<td>SafeSnapshot</td>
<td>Waiting for a snapshot for a READ ONLY DEFERRABLE transaction.</td>
<td></td>
</tr>
<tr>
<td>SyncRep</td>
<td>Waiting for confirmation from remote server during synchronous replication.</td>
<td></td>
</tr>
<tr>
<td>Timeout</td>
<td>BaseBackupThrottle</td>
<td>waiting during base backup when throttling activity.</td>
</tr>
<tr>
<td></td>
<td>PgSleep</td>
<td>waiting in process that called pg_sleep.</td>
</tr>
<tr>
<td></td>
<td>RecoveryApplyDelay</td>
<td>waiting to apply WAL at recovery because it is delayed.</td>
</tr>
<tr>
<td>IO</td>
<td>BufFileRead</td>
<td>waiting for a read from a buffered file.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>BufFileWrite</td>
<td></td>
<td>Waiting for a write to a buffered file.</td>
</tr>
<tr>
<td>ControlFileRead</td>
<td></td>
<td>Waiting for a read from the control file.</td>
</tr>
<tr>
<td>ControlFileSync</td>
<td></td>
<td>Waiting for the control file to reach stable storage.</td>
</tr>
<tr>
<td>ControlFileSyncUpdate</td>
<td></td>
<td>Waiting for an update to the control file to reach stable storage.</td>
</tr>
<tr>
<td>ControlFileWrite</td>
<td></td>
<td>Waiting for a write to the control file.</td>
</tr>
<tr>
<td>ControlFileWriteUpdate</td>
<td></td>
<td>Waiting for a write to update the control file.</td>
</tr>
<tr>
<td>CopyFileRead</td>
<td></td>
<td>Waiting for a read during a file copy operation.</td>
</tr>
<tr>
<td>CopyFileWrite</td>
<td></td>
<td>Waiting for a write during a file copy operation.</td>
</tr>
<tr>
<td>DataFileExtend</td>
<td></td>
<td>Waiting for a relation data file to be extended.</td>
</tr>
<tr>
<td>DataFileFlush</td>
<td></td>
<td>Waiting for a relation data file to reach stable storage.</td>
</tr>
<tr>
<td>DataFileImmediateSync</td>
<td></td>
<td>Waiting for an immediate synchronization of a relation data file to stable storage.</td>
</tr>
<tr>
<td>DataFilePrefetch</td>
<td></td>
<td>Waiting for an asynchronous prefetch from a relation data file.</td>
</tr>
<tr>
<td>DataFileRead</td>
<td></td>
<td>Waiting for a read from a relation data file.</td>
</tr>
<tr>
<td>DataFileSync</td>
<td></td>
<td>Waiting for changes to a relation data file to reach stable storage.</td>
</tr>
<tr>
<td>DataFileTruncate</td>
<td></td>
<td>Waiting for a relation data file to be truncated.</td>
</tr>
<tr>
<td>DataFileWrite</td>
<td></td>
<td>Waiting for a write to a relation data file.</td>
</tr>
<tr>
<td>DSMFillZeroWrite</td>
<td></td>
<td>Waiting to write zero bytes to a dynamic shared memory backing file.</td>
</tr>
<tr>
<td>LockFileAddToDataDirRead</td>
<td></td>
<td>Waiting for a read while adding a line to the data directory lock file.</td>
</tr>
<tr>
<td>LockFileAddToDataDirSync</td>
<td></td>
<td>Waiting for data to reach stable storage while adding a line to the data directory lock file.</td>
</tr>
<tr>
<td>LockFileAddToDataDirWrite</td>
<td></td>
<td>Waiting for a write while adding a line to the data directory lock file.</td>
</tr>
<tr>
<td>LockFileCreateRead</td>
<td></td>
<td>Waiting to read while creating the data directory lock file.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------------</td>
<td>--------------------------------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>LockFileCreateSync</td>
<td>LockFileCreateSync</td>
<td>Waiting for data to reach stable storage while creating the data directory lock file.</td>
</tr>
<tr>
<td>LockFileCreateWrite</td>
<td>LockFileCreateWrite</td>
<td>Waiting for a write while creating the data directory lock file.</td>
</tr>
<tr>
<td>LockFileReCheck-DataDirRead</td>
<td>LockFileReCheck-DataDirRead</td>
<td>Waiting for a read during recheck of the data directory lock file.</td>
</tr>
<tr>
<td>LogicalRewriteCheck-pointSync</td>
<td>LogicalRewriteCheck-pointSync</td>
<td>Waiting for logical rewrite mappings to reach stable storage during a checkpoint.</td>
</tr>
<tr>
<td>LogicalRewriteMappingSync</td>
<td>LogicalRewriteMappingSync</td>
<td>Waiting for mapping data to reach stable storage during a logical rewrite.</td>
</tr>
<tr>
<td>LogicalRewriteMappingWrite</td>
<td>LogicalRewriteMappingWrite</td>
<td>Waiting for a write of mapping data during a logical rewrite.</td>
</tr>
<tr>
<td>LogicalRewriteSync</td>
<td>LogicalRewriteSync</td>
<td>Waiting for logical rewrite mappings to reach stable storage.</td>
</tr>
<tr>
<td>LogicalRewriteWrite</td>
<td>LogicalRewriteWrite</td>
<td>Waiting for a write of logical rewrite mappings.</td>
</tr>
<tr>
<td>RelationMapRead</td>
<td>RelationMapRead</td>
<td>Waiting for a read of the relation map file.</td>
</tr>
<tr>
<td>RelationMapSync</td>
<td>RelationMapSync</td>
<td>Waiting for the relation map file to reach stable storage.</td>
</tr>
<tr>
<td>RelationMapWrite</td>
<td>RelationMapWrite</td>
<td>Waiting for a write to the relation map file.</td>
</tr>
<tr>
<td>ReorderBufferRead</td>
<td>ReorderBufferRead</td>
<td>Waiting for a read during reorder buffer management.</td>
</tr>
<tr>
<td>ReorderBufferWrite</td>
<td>ReorderBufferWrite</td>
<td>Waiting for a write during reorder buffer management.</td>
</tr>
<tr>
<td>ReorderLogicalMappingRead</td>
<td>ReorderLogicalMappingRead</td>
<td>Waiting for a read of a logical mapping during reorder buffer management.</td>
</tr>
<tr>
<td>ReplicationSlotRead</td>
<td>ReplicationSlotRead</td>
<td>Waiting for a read from a replication slot control file.</td>
</tr>
<tr>
<td>ReplicationSlotRestoreSync</td>
<td>ReplicationSlotRestoreSync</td>
<td>Waiting for a replication slot control file to reach stable storage while restoring it to memory.</td>
</tr>
<tr>
<td>ReplicationSlotSync</td>
<td>ReplicationSlotSync</td>
<td>Waiting for a replication slot control file to reach stable storage.</td>
</tr>
<tr>
<td>ReplicationSlotWrite</td>
<td>ReplicationSlotWrite</td>
<td>Waiting for a write to a replication slot control file.</td>
</tr>
<tr>
<td>SLRUFlushSync</td>
<td>SLRUFlushSync</td>
<td>Waiting for SLRU data to reach stable storage during a checkpoint or database shutdown.</td>
</tr>
<tr>
<td>SLRURead</td>
<td>SLRURead</td>
<td>Waiting for a read of an SLRU page.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>--------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>SLRUSync</td>
<td>Waiting for SLRU data to reach stable storage following a page write.</td>
<td></td>
</tr>
<tr>
<td>SLRUWrite</td>
<td>Waiting for a write of an SLRU page.</td>
<td></td>
</tr>
<tr>
<td>SnapbuildRead</td>
<td>Waiting for a read of a serialized historical catalog snapshot.</td>
<td></td>
</tr>
<tr>
<td>SnapbuildSync</td>
<td>Waiting for a serialized historical catalog snapshot to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>SnapbuildWrite</td>
<td>Waiting for a write of a serialized historical catalog snapshot.</td>
<td></td>
</tr>
<tr>
<td>TimelineHistoryFileSync</td>
<td>Waiting for a timeline history file received via streaming replication to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>TimelineHistoryFileWrite</td>
<td>Waiting for a write of a timeline history file received via streaming replication.</td>
<td></td>
</tr>
<tr>
<td>TimelineHistoryRead</td>
<td>Waiting for a read of a timeline history file.</td>
<td></td>
</tr>
<tr>
<td>TimelineHistorySync</td>
<td>Waiting for a newly created timeline history file to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>TimelineHistoryWrite</td>
<td>Waiting for a write of a newly created timeline history file.</td>
<td></td>
</tr>
<tr>
<td>TwophaseFileRead</td>
<td>Waiting for a read of a two phase state file.</td>
<td></td>
</tr>
<tr>
<td>TwophaseFileSync</td>
<td>Waiting for a two phase state file to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>TwophaseFileWrite</td>
<td>Waiting for a write of a two phase state file.</td>
<td></td>
</tr>
<tr>
<td>WALBootstrapSync</td>
<td>Waiting for WAL to reach stable storage during bootstrapping.</td>
<td></td>
</tr>
<tr>
<td>WALBootstrapWrite</td>
<td>Waiting for a write of a WAL page during bootstrapping.</td>
<td></td>
</tr>
<tr>
<td>WALCopyRead</td>
<td>Waiting for a read when creating a new WAL segment by copying an existing one.</td>
<td></td>
</tr>
<tr>
<td>WALCopySync</td>
<td>Waiting a new WAL segment created by copying an existing one to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>WALCopyWrite</td>
<td>Waiting for a write when creating a new WAL segment by copying an existing one.</td>
<td></td>
</tr>
<tr>
<td>WALInitSync</td>
<td>Waiting for a newly initialized WAL file to reach stable storage.</td>
<td></td>
</tr>
<tr>
<td>WALInitWrite</td>
<td>Waiting for a write while initializing a new WAL file.</td>
<td></td>
</tr>
</tbody>
</table>
### Wait Event Type

<table>
<thead>
<tr>
<th>Wait Event Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>WALRead</td>
<td>Waiting for a read from a WAL file.</td>
</tr>
<tr>
<td>WALSenderTimelineHistoryRead</td>
<td>Waiting for a read from a timeline history file during walsender timeline command.</td>
</tr>
<tr>
<td>WALSyncMethodAssign</td>
<td>Waiting for data to reach stable storage while assigning WAL sync method.</td>
</tr>
<tr>
<td>WALWrite</td>
<td>Waiting for a write to a WAL file.</td>
</tr>
</tbody>
</table>

**Note**

For tranches registered by extensions, the name is specified by extension and this will be displayed as `wait_event`. It is quite possible that user has registered the tranche in one of the backends (by having allocation in dynamic shared memory) in which case other backends won’t have that information, so we display `extension` for such cases.

Here is an example of how wait events can be viewed

```sql
SELECT pid, wait_event_type, wait_event FROM pg_stat_activity WHERE wait_event is NOT NULL;
```

<table>
<thead>
<tr>
<th>pid</th>
<th>wait_event_type</th>
<th>wait_event</th>
</tr>
</thead>
<tbody>
<tr>
<td>2540</td>
<td>Lock</td>
<td>relation</td>
</tr>
<tr>
<td>6644</td>
<td>LWLock</td>
<td>ProcArrayLock</td>
</tr>
</tbody>
</table>

(2 rows)

**Table 28.5. pg_stat_replication View**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of a WAL sender process</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>OID of the user logged into this WAL sender process</td>
</tr>
<tr>
<td>usename</td>
<td>name</td>
<td>Name of the user logged into this WAL sender process</td>
</tr>
<tr>
<td>application_name</td>
<td>text</td>
<td>Name of the application that is connected to this WAL sender</td>
</tr>
<tr>
<td>client_addr</td>
<td>inet</td>
<td>IP address of the client connected to this WAL sender. If this field is null, it indicates that the client is connected via a Unix socket on the server machine.</td>
</tr>
<tr>
<td>client_hostname</td>
<td>text</td>
<td>Host name of the connected client, as reported by a reverse DNS lookup of <code>client_addr</code>. This field will only be non-null for IP connections, and only when log_hostname is enabled.</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>-----------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>client_port</td>
<td>integer</td>
<td>TCP port number that the client is using for communication with this WAL sender, or -1 if a Unix socket is used</td>
</tr>
<tr>
<td>backend_start</td>
<td>timestamp with timezone</td>
<td>Time when this process was started, i.e., when the client connected to this WAL sender</td>
</tr>
<tr>
<td>backend_xmin</td>
<td>xid</td>
<td>This standby's xmin horizon reported by hot_standby_feedback.</td>
</tr>
<tr>
<td>state</td>
<td>text</td>
<td>Current WAL sender state. Possible values are:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• startup: This WAL sender is starting up.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• catchup: This WAL sender's connected standby is catching up with the primary.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• streaming: This WAL sender is streaming changes after its connected standby server has caught up with the primary.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• backup: This WAL sender is sending a backup.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• stopping: This WAL sender is stopping.</td>
</tr>
<tr>
<td>sent_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location sent on this connection</td>
</tr>
<tr>
<td>write_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location written to disk by this standby server</td>
</tr>
<tr>
<td>flush_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location flushed to disk by this standby server</td>
</tr>
<tr>
<td>replay_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location replayed into the database on this standby server</td>
</tr>
<tr>
<td>write_lag</td>
<td>interval</td>
<td>Time elapsed between flushing recent WAL locally and receiving notification that this standby server has written it (but not yet flushed it or applied it). This can be used to gauge the delay that synchronous_commit level remote_write incurred while committing if this server was configured as a synchronous standby.</td>
</tr>
</tbody>
</table>
## Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>flush_lag</td>
<td>interval</td>
<td>Time elapsed between flushing recent WAL locally and receiving notification that this standby server has written and flushed it (but not yet applied it). This can be used to gauge the delay that <code>synchronous_commit</code> level occurred while committing if this server was configured as a synchronous standby.</td>
</tr>
<tr>
<td>replay_lag</td>
<td>interval</td>
<td>Time elapsed between flushing recent WAL locally and receiving notification that this standby server has written, flushed and applied it. This can be used to gauge the delay that <code>synchronous_commit</code> level incurred while committing if this server was configured as a synchronous standby.</td>
</tr>
<tr>
<td>sync_priority</td>
<td>integer</td>
<td>Priority of this standby server for being chosen as the synchronous standby in a priority-based synchronous replication. This has no effect in a quorum-based synchronous replication.</td>
</tr>
<tr>
<td>sync_state</td>
<td>text</td>
<td>Synchronous state of this standby server. Possible values are:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• async: This standby server is asynchronous.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• potential: This standby server is now asynchronous, but can potentially become synchronous if one of current synchronous ones fails.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• sync: This standby server is synchronous.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• quorum: This standby server is considered a candidate for quorum standbys.</td>
</tr>
</tbody>
</table>

The `pg_stat_replication` view will contain one row per WAL sender process, showing statistics about replication to that sender's connected standby server. Only directly connected standbys are listed; no information is available about downstream standby servers.

The lag times reported in the `pg_stat_replication` view are measurements of the time taken for recent WAL to be written, flushed and replayed and for the sender to know about it. These times represent the commit delay that was (or would have been) introduced by each synchronous commit level, if the remote server was configured as a synchronous standby. For an asynchronous standby, the `replay_lag` column approximates the delay before recent transactions became visible to queries. If
the standby server has entirely caught up with the sending server and there is no more WAL activity, the most recently measured lag times will continue to be displayed for a short time and then show NULL.

Lag times work automatically for physical replication. Logical decoding plugins may optionally emit tracking messages; if they do not, the tracking mechanism will simply display NULL lag.

### Note

The reported lag times are not predictions of how long it will take for the standby to catch up with the sending server assuming the current rate of replay. Such a system would show similar times while new WAL is being generated, but would differ when the sender becomes idle. In particular, when the standby has caught up completely, `pg_stat_replication` shows the time taken to write, flush and replay the most recent reported WAL location rather than zero as some users might expect. This is consistent with the goal of measuring synchronous commit and transaction visibility delays for recent write transactions. To reduce confusion for users expecting a different model of lag, the lag columns revert to NULL after a short time on a fully replayed idle system. Monitoring systems should choose whether to represent this as missing data, zero or continue to display the last known value.

### Table 28.6. `pg_stat_wal_receiver` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of the WAL receiver process</td>
</tr>
<tr>
<td>status</td>
<td>text</td>
<td>Activity status of the WAL receiver process</td>
</tr>
<tr>
<td>receive_start_lsn</td>
<td><code>pg_lsn</code></td>
<td>First write-ahead log location used when WAL receiver is started</td>
</tr>
<tr>
<td>receive_start_tli</td>
<td>integer</td>
<td>First timeline number used when WAL receiver is started</td>
</tr>
<tr>
<td>received_lsn</td>
<td><code>pg_lsn</code></td>
<td>Last write-ahead log location already received and flushed to disk, the initial value of this field being the first log location used when WAL receiver is started</td>
</tr>
<tr>
<td>received_tli</td>
<td>integer</td>
<td>Timeline number of last write-ahead log location received and flushed to disk, the initial value of this field being the timeline number of the first log location used when WAL receiver is started</td>
</tr>
<tr>
<td>last_msg_send_time</td>
<td><code>timestamp with time zone</code></td>
<td>Send time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>last_msg_receipt_time</td>
<td><code>timestamp with time zone</code></td>
<td>Receipt time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>latest_end_lsn</td>
<td><code>pg_lsn</code></td>
<td>Last write-ahead log location reported to origin WAL sender</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------</td>
<td>-----------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>latest_end_time</td>
<td>timestamp with time zone</td>
<td>Time of last write-ahead log location reported to origin WAL sender</td>
</tr>
<tr>
<td>slot_name</td>
<td>text</td>
<td>Replication slot name used by this WAL receiver</td>
</tr>
<tr>
<td>sender_host</td>
<td>text</td>
<td>Host of the PostgreSQL instance this WAL receiver is connected to. This can be a host name, an IP address, or a directory path if the connection is via Unix socket. (The path case can be distinguished because it will always be an absolute path, beginning with <code>/</code>.)</td>
</tr>
<tr>
<td>sender_port</td>
<td>integer</td>
<td>Port number of the PostgreSQL instance this WAL receiver is connected to.</td>
</tr>
<tr>
<td>conninfo</td>
<td>text</td>
<td>Connection string used by this WAL receiver, with security-sensitive fields obfuscated.</td>
</tr>
</tbody>
</table>

The `pg_stat_wal_receiver` view will contain only one row, showing statistics about the WAL receiver from that receiver's connected server.

### Table 28.7. `pg_stat_subscription` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>subid</td>
<td>oid</td>
<td>OID of the subscription</td>
</tr>
<tr>
<td>subname</td>
<td>text</td>
<td>Name of the subscription</td>
</tr>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of the subscription worker process</td>
</tr>
<tr>
<td>relid</td>
<td>Oid</td>
<td>OID of the relation that the worker is synchronizing; null for the main apply worker</td>
</tr>
<tr>
<td>received_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location received, the initial value of this field being 0</td>
</tr>
<tr>
<td>last_msg_send_time</td>
<td>timestamp with time zone</td>
<td>Send time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>last_msg_receipt_time</td>
<td>timestamp with time zone</td>
<td>Receipt time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>latest_end_lsn</td>
<td>pg_lsn</td>
<td>Last write-ahead log location reported to origin WAL sender</td>
</tr>
<tr>
<td>latest_end_time</td>
<td>timestamp with time zone</td>
<td>Time of last write-ahead log location reported to origin WAL sender</td>
</tr>
</tbody>
</table>

The `pg_stat_subscription` view will contain one row per subscription for main worker (with null PID if the worker is not running), and additional rows for workers handling the initial data copy of the subscribed tables.
Table 28.8. pg_stat_ssl View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of a backend or WAL sender process</td>
</tr>
<tr>
<td>ssl</td>
<td>boolean</td>
<td>True if SSL is used on this connection</td>
</tr>
<tr>
<td>version</td>
<td>text</td>
<td>Version of SSL in use, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>cipher</td>
<td>text</td>
<td>Name of SSL cipher in use, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>bits</td>
<td>integer</td>
<td>Number of bits in the encryption algorithm used, or NULL if SSL is not used on this connection</td>
</tr>
<tr>
<td>compression</td>
<td>boolean</td>
<td>True if SSL compression is in use, false if not, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>clientdn</td>
<td>text</td>
<td>Distinguished Name (DN) field from the client certificate used, or NULL if no client certificate was supplied or if SSL is not in use on this connection. This field is truncated if the DN field is longer than NAMEDATALEN (64 characters in a standard build)</td>
</tr>
</tbody>
</table>

The pg_stat_ssl view will contain one row per backend or WAL sender process, showing statistics about SSL usage on this connection. It can be joined to pg_stat_activity or pg_stat_replication on the pid column to get more details about the connection.

Table 28.9. pg_stat_archiver View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>archived_count</td>
<td>bigint</td>
<td>Number of WAL files that have been successfully archived</td>
</tr>
<tr>
<td>last_archived_wal</td>
<td>text</td>
<td>Name of the last WAL file successfully archived</td>
</tr>
<tr>
<td>last_archived_time</td>
<td>timestamp with time zone</td>
<td>Time of the last successful archive operation</td>
</tr>
<tr>
<td>failed_count</td>
<td>bigint</td>
<td>Number of failed attempts for archiving WAL files</td>
</tr>
<tr>
<td>last_failed_wal</td>
<td>text</td>
<td>Name of the WAL file of the last failed archival operation</td>
</tr>
<tr>
<td>last_failed_time</td>
<td>timestamp with time zone</td>
<td>Time of the last failed archival operation</td>
</tr>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset</td>
</tr>
</tbody>
</table>

The pg_stat_archiver view will always have a single row, containing data about the archiver process of the cluster.
### Table 28.10. pg_stat_bgwriter View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>checkpoints_timed</td>
<td>bigint</td>
<td>Number of scheduled checkpoints that have been performed</td>
</tr>
<tr>
<td>checkpoints_req</td>
<td>bigint</td>
<td>Number of requested checkpoints that have been performed</td>
</tr>
<tr>
<td>checkpoint_write_time</td>
<td>double precision</td>
<td>Total amount of time that has been spent in the portion of checkpoint processing where files are written to disk, in milliseconds</td>
</tr>
<tr>
<td>checkpoint_sync_time</td>
<td>double precision</td>
<td>Total amount of time that has been spent in the portion of checkpoint processing where files are synchronized to disk, in milliseconds</td>
</tr>
<tr>
<td>buffers_checkpoint</td>
<td>bigint</td>
<td>Number of buffers written during checkpoints</td>
</tr>
<tr>
<td>buffers_clean</td>
<td>bigint</td>
<td>Number of buffers written by the background writer</td>
</tr>
<tr>
<td>maxwritten_clean</td>
<td>bigint</td>
<td>Number of times the background writer stopped a cleaning scan because it had written too many buffers</td>
</tr>
<tr>
<td>buffers_backend</td>
<td>bigint</td>
<td>Number of buffers written directly by a backend</td>
</tr>
<tr>
<td>buffers_backend_fsync</td>
<td>bigint</td>
<td>Number of times a backend had to execute its own fsync call (normally the background writer handles those even when the backend does its own write)</td>
</tr>
<tr>
<td>buffers_alloc</td>
<td>bigint</td>
<td>Number of buffers allocated</td>
</tr>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset</td>
</tr>
</tbody>
</table>

The `pg_stat_bgwriter` view will always have a single row, containing global data for the cluster.

### Table 28.11. pg_stat_database View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of a database</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of this database</td>
</tr>
<tr>
<td>numbackends</td>
<td>integer</td>
<td>Number of backends currently connected to this database. This is the only column in this view that returns a value reflecting current state; all other columns return the accumulated values since the last reset.</td>
</tr>
<tr>
<td>xact_commit</td>
<td>bigint</td>
<td>Number of transactions in this database that have been committed</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>---------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>xact_rollback</td>
<td>bigint</td>
<td>Number of transactions in this database that have been rolled back</td>
</tr>
<tr>
<td>blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read in this database</td>
</tr>
<tr>
<td>blks_hit</td>
<td>bigint</td>
<td>Number of times disk blocks were found already in the buffer cache, so that a read was not necessary (this only includes hits in the PostgreSQL buffer cache, not the operating system's file system cache)</td>
</tr>
<tr>
<td>tup Returned</td>
<td>bigint</td>
<td>Number of rows returned by queries in this database</td>
</tr>
<tr>
<td>tup_fetched</td>
<td>bigint</td>
<td>Number of rows fetched by queries in this database</td>
</tr>
<tr>
<td>tup_inserted</td>
<td>bigint</td>
<td>Number of rows inserted by queries in this database</td>
</tr>
<tr>
<td>tup_updated</td>
<td>bigint</td>
<td>Number of rows updated by queries in this database</td>
</tr>
<tr>
<td>tup_deleted</td>
<td>bigint</td>
<td>Number of rows deleted by queries in this database</td>
</tr>
<tr>
<td>conflicts</td>
<td>bigint</td>
<td>Number of queries canceled due to conflicts with recovery in this database. (Conflicts occur only on standby servers; see pg_stat_database_conflicts for details.)</td>
</tr>
<tr>
<td>temp_files</td>
<td>bigint</td>
<td>Number of temporary files created by queries in this database. All temporary files are counted, regardless of why the temporary file was created (e.g., sorting or hashing), and regardless of the log_temp_files setting.</td>
</tr>
<tr>
<td>temp_bytes</td>
<td>bigint</td>
<td>Total amount of data written to temporary files by queries in this database. All temporary files are counted, regardless of why the temporary file was created, and regardless of the log_temp_files setting.</td>
</tr>
<tr>
<td>deadlocks</td>
<td>bigint</td>
<td>Number of deadlocks detected in this database</td>
</tr>
<tr>
<td>blk_read_time</td>
<td>double precision</td>
<td>Time spent reading data file blocks by backends in this database, in milliseconds</td>
</tr>
<tr>
<td>blk_write_time</td>
<td>double precision</td>
<td>Time spent writing data file blocks by backends in this database, in milliseconds</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset</td>
</tr>
</tbody>
</table>

The `pg_stat_database` view will contain one row for each database in the cluster, showing database-wide statistics.

### Table 28.12. `pg_stat_database_conflicts` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of a database</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of this database</td>
</tr>
<tr>
<td>confl_tablespace</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to dropped table spaces</td>
</tr>
<tr>
<td>confl_lock</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to lock timeouts</td>
</tr>
<tr>
<td>confl_snapshot</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to old snapshots</td>
</tr>
<tr>
<td>confl_bufferpin</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to pinned buffers</td>
</tr>
<tr>
<td>confl_deadlock</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to deadlocks</td>
</tr>
</tbody>
</table>

The `pg_stat_database_conflicts` view will contain one row per database, showing database-wide statistics about query cancels occurring due to conflicts with recovery on standby servers. This view will only contain information on standby servers, since conflicts do not occur on master servers.

### Table 28.13. `pg_stat_all_tables` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of a table</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema that this table is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of this table</td>
</tr>
<tr>
<td>seq_scan</td>
<td>bigint</td>
<td>Number of sequential scans initiated on this table</td>
</tr>
<tr>
<td>seq_tup_read</td>
<td>bigint</td>
<td>Number of live rows fetched by sequential scans</td>
</tr>
<tr>
<td>idx_scan</td>
<td>bigint</td>
<td>Number of index scans initiated on this table</td>
</tr>
<tr>
<td>idx_tup_fetch</td>
<td>bigint</td>
<td>Number of live rows fetched by index scans</td>
</tr>
<tr>
<td>n_tup_ins</td>
<td>bigint</td>
<td>Number of rows inserted</td>
</tr>
<tr>
<td>n_tup_upd</td>
<td>bigint</td>
<td>Number of rows updated (includes HOT updated rows)</td>
</tr>
<tr>
<td>n_tup_del</td>
<td>bigint</td>
<td>Number of rows deleted</td>
</tr>
</tbody>
</table>
### Monitoring Database Activity

#### Table 28.14. `pg_stat_all_tables` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>n_tup_hot_upd</code></td>
<td>bigint</td>
<td>Number of rows HOT updated (i.e., with no separate index update required)</td>
</tr>
<tr>
<td><code>n_live_tup</code></td>
<td>bigint</td>
<td>Estimated number of live rows</td>
</tr>
<tr>
<td><code>n_dead_tup</code></td>
<td>bigint</td>
<td>Estimated number of dead rows</td>
</tr>
<tr>
<td><code>n_mod_since_analyze</code></td>
<td>bigint</td>
<td>Estimated number of rows modified since this table was last analyzed</td>
</tr>
<tr>
<td><code>last_vacuum</code></td>
<td>timestamp with time zone</td>
<td>Last time at which this table was manually vacuumed (not counting VACUUM FULL)</td>
</tr>
<tr>
<td><code>last_autovacuum</code></td>
<td>timestamp with time zone</td>
<td>Last time at which this table was vacuumed by the autovacuum daemon</td>
</tr>
<tr>
<td><code>last_analyze</code></td>
<td>timestamp with time zone</td>
<td>Last time at which this table was manually analyzed</td>
</tr>
<tr>
<td><code>last_autoanalyze</code></td>
<td>timestamp with time zone</td>
<td>Last time at which this table was analyzed by the autovacuum daemon</td>
</tr>
<tr>
<td><code>vacuum_count</code></td>
<td>bigint</td>
<td>Number of times this table has been manually vacuumed (not counting VACUUM FULL)</td>
</tr>
<tr>
<td><code>autovacuum_count</code></td>
<td>bigint</td>
<td>Number of times this table has been vacuumed by the autovacuum daemon</td>
</tr>
<tr>
<td><code>analyze_count</code></td>
<td>bigint</td>
<td>Number of times this table has been manually analyzed</td>
</tr>
<tr>
<td><code>autoanalyze_count</code></td>
<td>bigint</td>
<td>Number of times this table has been analyzed by the autovacuum daemon</td>
</tr>
</tbody>
</table>

The `pg_stat_all_tables` view will contain one row for each table in the current database (including TOAST tables), showing statistics about accesses to that specific table. The `pg_stat_user_tables` and `pg_stat_sys_tables` views contain the same information, but filtered to only show user and system tables respectively.

#### Table 28.14. `pg_stat_all_indexes` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>relid</code></td>
<td>oid</td>
<td>OID of the table for this index</td>
</tr>
<tr>
<td><code>indexrelid</code></td>
<td>oid</td>
<td>OID of this index</td>
</tr>
<tr>
<td><code>schemaname</code></td>
<td>name</td>
<td>Name of the schema this index is in</td>
</tr>
<tr>
<td><code>relname</code></td>
<td>name</td>
<td>Name of the table for this index</td>
</tr>
<tr>
<td><code>indexrelname</code></td>
<td>name</td>
<td>Name of this index</td>
</tr>
<tr>
<td><code>idx_scan</code></td>
<td>bigint</td>
<td>Number of index scans initiated on this index</td>
</tr>
<tr>
<td><code>idx_tup_read</code></td>
<td>bigint</td>
<td>Number of index entries returned by scans on this index</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>idx_tup_fetch</td>
<td>bigint</td>
<td>Number of live table rows fetched by simple index scans using this index</td>
</tr>
</tbody>
</table>

The `pg_stat_all_indexes` view will contain one row for each index in the current database, showing statistics about accesses to that specific index. The `pg_stat_user_indexes` and `pg_stat_sys_indexes` views contain the same information, but filtered to only show user and system indexes respectively.

Indexes can be used by simple index scans, “bitmap” index scans, and the optimizer. In a bitmap scan the output of several indexes can be combined via AND or OR rules, so it is difficult to associate individual heap row fetches with specific indexes when a bitmap scan is used. Therefore, a bitmap scan increments the `pg_stat_all_indexes.idx_tup_read` count(s) for the index(es) it uses, and it increments the `pg_stat_all_tables.idx_tup_fetch` count for the table, but it does not affect `pg_stat_all_indexes.idx_tup_fetch`. The optimizer also accesses indexes to check for supplied constants whose values are outside the recorded range of the optimizer statistics because the optimizer statistics might be stale.

### Note

The `idx_tup_read` and `idx_tup_fetch` counts can be different even without any use of bitmap scans, because `idx_tup_read` counts index entries retrieved from the index while `idx_tup_fetch` counts live rows fetched from the table. The latter will be less if any dead or not-yet-committed rows are fetched using the index, or if any heap fetches are avoided by means of an index-only scan.

Table 28.15. `pg_statio_all_tables` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of a table</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema that this table is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of this table</td>
</tr>
<tr>
<td>heap_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this table</td>
</tr>
<tr>
<td>heap_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this table</td>
</tr>
<tr>
<td>idx_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from all indexes on this table</td>
</tr>
<tr>
<td>idx_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in all indexes on this table</td>
</tr>
<tr>
<td>toast_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this table's TOAST table (if any)</td>
</tr>
<tr>
<td>toast_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this table's TOAST table (if any)</td>
</tr>
<tr>
<td>tidx_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this table's TOAST table indexes (if any)</td>
</tr>
<tr>
<td>tidx_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this table's TOAST table indexes (if any)</td>
</tr>
</tbody>
</table>
Monitoring Database Activity

The `pg_statio_all_tables` view will contain one row for each table in the current database (including TOAST tables), showing statistics about I/O on that specific table. The `pg_statio_user_tables` and `pg_statio_sys_tables` views contain the same information, but filtered to only show user and system tables respectively.

**Table 28.16. pg_statio_all_indexes View**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of the table for this index</td>
</tr>
<tr>
<td>indexrelid</td>
<td>oid</td>
<td>OID of this index</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this index is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of the table for this index</td>
</tr>
<tr>
<td>indexrelname</td>
<td>name</td>
<td>Name of this index</td>
</tr>
<tr>
<td>idx_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this index</td>
</tr>
<tr>
<td>idx_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this index</td>
</tr>
</tbody>
</table>

The `pg_statio_all_indexes` view will contain one row for each index in the current database, showing statistics about I/O on that specific index. The `pg_statio_user_indexes` and `pg_statio_sys_indexes` views contain the same information, but filtered to only show user and system indexes respectively.

**Table 28.17. pg_statio_all_sequences View**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of a sequence</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this sequence is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of this sequence</td>
</tr>
<tr>
<td>blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this sequence</td>
</tr>
<tr>
<td>blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this sequence</td>
</tr>
</tbody>
</table>

The `pg_statio_all_sequences` view will contain one row for each sequence in the current database, showing statistics about I/O on that specific sequence.

**Table 28.18. pg_stat_user_functions View**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>funcid</td>
<td>oid</td>
<td>OID of a function</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this function is in</td>
</tr>
<tr>
<td>funcname</td>
<td>name</td>
<td>Name of this function</td>
</tr>
<tr>
<td>calls</td>
<td>bigint</td>
<td>Number of times this function has been called</td>
</tr>
<tr>
<td>total_time</td>
<td>double precision</td>
<td>Total time spent in this function and all other functions called by it, in milliseconds</td>
</tr>
</tbody>
</table>

723
Monitoring Database Activity

### Column Types and Descriptions

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>self_time</td>
<td>double precision</td>
<td>Total time spent in this function itself, not including other functions called by it, in milliseconds</td>
</tr>
</tbody>
</table>

The `pg_stat_user_functions` view will contain one row for each tracked function, showing statistics about executions of that function. The `track_functions` parameter controls exactly which functions are tracked.

### 28.2.3. Statistics Functions

Other ways of looking at the statistics can be set up by writing queries that use the same underlying statistics access functions used by the standard views shown above. For details such as the functions’ names, consult the definitions of the standard views. (For example, in `psql` you could issue `\d+ pg_stat_activity`.) The access functions for per-database statistics take a database OID as an argument to identify which database to report on. The per-table and per-index functions take a table or index OID. The functions for per-function statistics take a function OID. Note that only tables, indexes, and functions in the current database can be seen with these functions.

Additional functions related to statistics collection are listed in Table 28.19.

#### Table 28.19. Additional Statistics Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_backend_pid()</code></td>
<td>integer</td>
<td>Process ID of the server process handling the current session</td>
</tr>
<tr>
<td><code>pg_stat_get_activity(integer)</code></td>
<td>setof record</td>
<td>Returns a record of information about the backend with the specified PID, or one record for each active backend in the system if NULL is specified. The fields returned are a subset of those in the <code>pg_stat_activity</code> view.</td>
</tr>
<tr>
<td><code>pg_stat_get_snapshot_timestamp()</code></td>
<td>timestamp with time zone</td>
<td>Returns the timestamp of the current statistics snapshot</td>
</tr>
<tr>
<td><code>pg_stat_clear_snapshot()</code></td>
<td>void</td>
<td>Discard the current statistics snapshot</td>
</tr>
<tr>
<td><code>pg_stat_reset()</code></td>
<td>void</td>
<td>Reset all statistics counters for the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others.)</td>
</tr>
<tr>
<td><code>pg_stat_reset_shared(text)</code></td>
<td>void</td>
<td>Reset some cluster-wide statistics counters to zero, depending on the argument (requires superuser privileges by default, but EXECUTE for this function can be granted to others). Calling <code>pg_stat_reset_shared('bgwriter')</code> will zero all the counters shown in the <code>pg_stat_bgelog</code> view. Calling <code>pg_stat_reset_shared('archive')</code> will zero all the counters shown in the <code>pg_stat_archiver</code> view.</td>
</tr>
</tbody>
</table>

724
### Monitoring Database Activity

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_reset_single_table_counts(oid)</td>
<td>void</td>
<td>Reset statistics for a single table or index in the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others)</td>
</tr>
<tr>
<td>pg_stat_reset_single_function_counts(oid)</td>
<td>void</td>
<td>Reset statistics for a single function in the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others)</td>
</tr>
</tbody>
</table>

`pg_stat_get_activity`, the underlying function of the `pg_stat_activity` view, returns a set of records containing all the available information about each backend process. Sometimes it may be more convenient to obtain just a subset of this information. In such cases, an older set of per-backend statistics access functions can be used; these are shown in Table 28.20. These access functions use a backend ID number, which ranges from one to the number of currently active backends. The function `pg_stat_get_backend_idset` provides a convenient way to generate one row for each active backend for invoking these functions. For example, to show the PIDs and current queries of all backends:

```sql
SELECT pg_stat_get_backend_pid(s.backendid) AS pid,
       pg_stat_get_backend_activity(s.backendid) AS query
FROM (SELECT pg_stat_get_backend_idset() AS backendid) AS s;
```

### Table 28.20. Per-Backend Statistics Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_get_backend_idset()</td>
<td>setof integer</td>
<td>Set of currently active backend ID numbers (from 1 to the number of active backends)</td>
</tr>
<tr>
<td>pg_stat_get_backend_activity(integer)</td>
<td>text</td>
<td>Text of this backend's most recent query</td>
</tr>
<tr>
<td>pg_stat_get_backend_activity_s_start(integer)</td>
<td>timestamp with time zone</td>
<td>Time when the most recent query was started</td>
</tr>
<tr>
<td>pg_stat_get_backend_client_addr(integer)</td>
<td>inet</td>
<td>IP address of the client connected to this backend</td>
</tr>
<tr>
<td>pg_stat_get_backend_client_port(integer)</td>
<td>integer</td>
<td>TCP port number that the client is using for communication</td>
</tr>
<tr>
<td>pg_stat_get_backend_dbid(integer)</td>
<td>oid</td>
<td>OID of the database this backend is connected to</td>
</tr>
<tr>
<td>pg_stat_get_backend_pid(integer)</td>
<td>integer</td>
<td>Process ID of this backend</td>
</tr>
<tr>
<td>pg_stat_get_backend_start(integer)</td>
<td>timestamp with timezone</td>
<td>Time when this process was started</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------------------------------</td>
<td>-------------</td>
<td>-------------------------------------------------------</td>
</tr>
<tr>
<td>pg_stat_get_backend_userid(integer)</td>
<td>oid</td>
<td>OID of the user logged into this backend</td>
</tr>
<tr>
<td>pg_stat_get_backend_wait_event_type(integer)</td>
<td>text</td>
<td>Wait event type name if backend is currently waiting, otherwise NULL. See Table 28.4 for details.</td>
</tr>
<tr>
<td>pg_stat_get_backend_wait_event(integer)</td>
<td>text</td>
<td>Wait event name if backend is currently waiting, otherwise NULL. See Table 28.4 for details.</td>
</tr>
<tr>
<td>pg_stat_get_backend_xact_start(integer)</td>
<td>timestamp with timezone</td>
<td>Time when the current transaction was started</td>
</tr>
</tbody>
</table>

### 28.3. Viewing Locks

Another useful tool for monitoring database activity is the `pg_locks` system table. It allows the database administrator to view information about the outstanding locks in the lock manager. For example, this capability can be used to:

- View all the locks currently outstanding, all the locks on relations in a particular database, all the locks on a particular relation, or all the locks held by a particular PostgreSQL session.
- Determine the relation in the current database with the most ungranted locks (which might be a source of contention among database clients).
- Determine the effect of lock contention on overall database performance, as well as the extent to which contention varies with overall database traffic.

Details of the `pg_locks` view appear in Section 52.73. For more information on locking and managing concurrency with PostgreSQL, refer to Chapter 13.

### 28.4. Progress Reporting

PostgreSQL has the ability to report the progress of certain commands during command execution. Currently, the only command which supports progress reporting is `VACUUM`. This may be expanded in the future.

#### 28.4.1. VACUUM Progress Reporting

Whenever `VACUUM` is running, the `pg_stat_progress_vacuum` view will contain one row for each backend (including autovacuum worker processes) that is currently vacuuming. The tables below describe the information that will be reported and provide information about how to interpret it. Progress reporting is not currently supported for `VACUUM FULL` and backends running `VACUUM FULL` will not be listed in this view.

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of backend.</td>
</tr>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of the database to which this backend is connected.</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of the database to which this backend is connected.</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>---------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of the table being vacuumed.</td>
</tr>
<tr>
<td>phase</td>
<td>text</td>
<td>Current processing phase of vacuum. See Table 28.22.</td>
</tr>
<tr>
<td>heap_blks_total</td>
<td>bigint</td>
<td>Total number of heap blocks in the table. This number is reported as of the beginning of the scan; blocks added later will not be (and need not be) visited by this VACUUM.</td>
</tr>
<tr>
<td>heap_blks_scanned</td>
<td>bigint</td>
<td>Number of heap blocks scanned. Because the visibility map is used to optimize scans, some blocks will be skipped without inspection; skipped blocks are included in this total, so that this number will eventually become equal to heap_blks_total when the vacuum is complete. This counter only advances when the phase is scanning heap.</td>
</tr>
<tr>
<td>heap_blks_vacuumed</td>
<td>bigint</td>
<td>Number of heap blocks vacuumed. Unless the table has no indexes, this counter only advances when the phase is vacuuming heap. Blocks that contain no dead tuples are skipped, so the counter may sometimes skip forward in large increments.</td>
</tr>
<tr>
<td>index_vacuum_count</td>
<td>bigint</td>
<td>Number of completed index vacuum cycles.</td>
</tr>
<tr>
<td>max_dead_tuples</td>
<td>bigint</td>
<td>Number of dead tuples that we can store before needing to perform an index vacuum cycle, based on maintenance_work_mem.</td>
</tr>
<tr>
<td>num_dead_tuples</td>
<td>bigint</td>
<td>Number of dead tuples collected since the last index vacuum cycle.</td>
</tr>
</tbody>
</table>

Table 28.22. VACUUM phases

<table>
<thead>
<tr>
<th>Phase</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>initializing</td>
<td>VACUUM is preparing to begin scanning the heap. This phase is expected to be very brief.</td>
</tr>
<tr>
<td>scanning heap</td>
<td>VACUUM is currently scanning the heap. It will prune and defragment each page if required, and possibly perform freezing activity. The heap_blks_scanned column can be used to monitor the progress of the scan.</td>
</tr>
<tr>
<td>vacuuming indexes</td>
<td>VACUUM is currently vacuuming the indexes. If a table has any indexes, this will happen at least</td>
</tr>
</tbody>
</table>
### Monitoring Database Activity

<table>
<thead>
<tr>
<th>Phase</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>once per vacuum, after the heap has been completely scanned. It may happen multiple times per vacuum if maintenance_work_mem is insufficient to store the number of dead tuples found.</td>
<td></td>
</tr>
<tr>
<td>vacuuming heap</td>
<td>VACUUM is currently vacuuming the heap. Vacuuming the heap is distinct from scanning the heap, and occurs after each instance of vacuuming indexes. If heap_blks_scanned is less than heap_blks_total, the system will return to scanning the heap after this phase is completed; otherwise, it will begin cleaning up indexes after this phase is completed.</td>
</tr>
<tr>
<td>cleaning up indexes</td>
<td>VACUUM is currently cleaning up indexes. This occurs after the heap has been completely scanned and all vacuuming of the indexes and the heap has been completed.</td>
</tr>
<tr>
<td>truncating heap</td>
<td>VACUUM is currently truncating the heap so as to return empty pages at the end of the relation to the operating system. This occurs after cleaning up indexes.</td>
</tr>
<tr>
<td>performing final cleanup</td>
<td>VACUUM is performing final cleanup. During this phase, VACUUM will vacuum the free space map, update statistics in pg_class, and report statistics to the statistics collector. When this phase is completed, VACUUM will end.</td>
</tr>
</tbody>
</table>

### 28.5. Dynamic Tracing

PostgreSQL provides facilities to support dynamic tracing of the database server. This allows an external utility to be called at specific points in the code and thereby trace execution.

A number of probes or trace points are already inserted into the source code. These probes are intended to be used by database developers and administrators. By default the probes are not compiled into PostgreSQL; the user needs to explicitly tell the configure script to make the probes available.

Currently, the DTrace\(^1\) utility is supported, which, at the time of this writing, is available on Solaris, macOS, FreeBSD, NetBSD, and Oracle Linux. The SystemTap\(^2\) project for Linux provides a DTrace equivalent and can also be used. Supporting other dynamic tracing utilities is theoretically possible by changing the definitions for the macros in src/include/utils/probes.h.

#### 28.5.1. Compiling for Dynamic Tracing

By default, probes are not available, so you will need to explicitly tell the configure script to make the probes available in PostgreSQL. To include DTrace support specify `--enable-dtrace` to configure. See Section 16.4 for further information.

#### 28.5.2. Built-in Probes

A number of standard probes are provided in the source code, as shown in Table 28.23; Table 28.24 shows the types used in the probes. More probes can certainly be added to enhance PostgreSQL’s observability.

---

\(^1\) [https://en.wikipedia.org/wiki/DTrace](https://en.wikipedia.org/wiki/DTrace)

\(^2\) [http://sourceware.org/systemtap/](http://sourceware.org/systemtap/)
### Table 28.23. Built-in DTrace Probes

<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>transaction-start</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires at the start of a new transaction. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>transaction-commit</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires when a transaction completes successfully. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>transaction-abort</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires when a transaction completes unsuccessfully. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>query-start</td>
<td>(const char *)</td>
<td>Probe that fires when the processing of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-done</td>
<td>(const char *)</td>
<td>Probe that fires when the processing of a query is complete. arg0 is the query string.</td>
</tr>
<tr>
<td>query-parse-start</td>
<td>(const char *)</td>
<td>Probe that fires when the parsing of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-parse-done</td>
<td>(const char *)</td>
<td>Probe that fires when the parsing of a query is complete. arg0 is the query string.</td>
</tr>
<tr>
<td>query-rewrite-start</td>
<td>(const char *)</td>
<td>Probe that fires when the rewriting of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-rewrite-done</td>
<td>(const char *)</td>
<td>Probe that fires when the rewriting of a query is complete. arg0 is the query string.</td>
</tr>
<tr>
<td>query-plan-start</td>
<td>()</td>
<td>Probe that fires when the planning of a query is started.</td>
</tr>
<tr>
<td>query-plan-done</td>
<td>()</td>
<td>Probe that fires when the planning of a query is complete.</td>
</tr>
<tr>
<td>query-execute-start</td>
<td>()</td>
<td>Probe that fires when the execution of a query is started.</td>
</tr>
<tr>
<td>query-execute-done</td>
<td>()</td>
<td>Probe that fires when the execution of a query is complete.</td>
</tr>
<tr>
<td>statement-status</td>
<td>(const char *)</td>
<td>Probe that fires anytime the server process updates its <code>pg_stat_activity.s-status</code>. arg0 is the new status string.</td>
</tr>
<tr>
<td>checkpoint-start</td>
<td>(int)</td>
<td>Probe that fires when a checkpoint is started. arg0 holds the bitwise flags used to distinguish different checkpoint types, such as shutdown, immediate or force.</td>
</tr>
<tr>
<td>checkpoint-done</td>
<td>(int, int, int, int, int)</td>
<td>Probe that fires when a checkpoint is complete. (The probes listed next fire in sequence dur-</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------</td>
<td>------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>clog-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the CLOG portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>clog-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the CLOG portion of a checkpoint is complete. arg0 has the same meaning as for clog-checkpoint-start.</td>
</tr>
<tr>
<td>subtrans-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the SUBTRANS portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>subtrans-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the SUBTRANS portion of a checkpoint is complete. arg0 has the same meaning as for subtrans-checkpoint-start.</td>
</tr>
<tr>
<td>multixact-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the MultiXact portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>multixact-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the MultiXact portion of a checkpoint is complete. arg0 has the same meaning as for multixact-checkpoint-start.</td>
</tr>
<tr>
<td>buffer-checkpoint-start</td>
<td>(int)</td>
<td>Probe that fires when the buffer-writing portion of a checkpoint is started. arg0 holds the bitwise flags used to distinguish different checkpoint types, such as shutdown, immediate or force.</td>
</tr>
<tr>
<td>buffer-sync-start</td>
<td>(int, int)</td>
<td>Probe that fires when we begin to write dirty buffers during checkpoint (after identifying which buffers must be written). arg0 is the total number of buffers. arg1 is the number that are currently dirty and need to be written.</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------------------</td>
<td>---------------------</td>
<td>--------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>buffer-sync-written</td>
<td>(int)</td>
<td>Probe that fires after each buffer is written during checkpoint. arg0 is the ID number of the buffer.</td>
</tr>
<tr>
<td>buffer-sync-done</td>
<td>(int, int, int)</td>
<td>Probe that fires when all dirty buffers have been written. arg0 is the total number of buffers. arg1 is the number of buffers actually written by the checkpoint process. arg2 is the number that were expected to be written (arg1 of buffer-sync-start); any difference reflects other processes flushing buffers during the checkpoint.</td>
</tr>
<tr>
<td>buffer-checkpoint-sync-start</td>
<td>()</td>
<td>Probe that fires after dirty buffers have been written to the kernel, and before starting to issue fsync requests.</td>
</tr>
<tr>
<td>buffer-checkpoint-done</td>
<td>()</td>
<td>Probe that fires when syncing of buffers to disk is complete.</td>
</tr>
<tr>
<td>twophase-checkpoint-start</td>
<td>()</td>
<td>Probe that fires when the two-phase portion of a checkpoint is started.</td>
</tr>
<tr>
<td>twophase-checkpoint-done</td>
<td>()</td>
<td>Probe that fires when the two-phase portion of a checkpoint is complete.</td>
</tr>
<tr>
<td>buffer-read-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, bool)</td>
<td>Probe that fires when a buffer read is started. arg0 and arg1 contain the fork and block numbers of the page (but arg1 will be -1 if this is a relation extension request). arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer. arg6 is true for a relation extension request, false for normal read.</td>
</tr>
<tr>
<td>buffer-read-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, bool, bool)</td>
<td>Probe that fires when a buffer read is complete. arg0 and arg1 contain the fork and block numbers of the page (if this is a relation extension request, arg1 now contains the block number of the newly added block). arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the back-</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------</td>
<td>------------------------------------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>buffer-flush-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires before issuing any write request for a shared buffer. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation.</td>
</tr>
<tr>
<td>buffer-flush-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a write request is complete. (Note that this just reflects the time to pass the data to the kernel; it’s typically not actually been written to disk yet.) The arguments are the same as for buffer-flush-start.</td>
</tr>
<tr>
<td>buffer-write-dirty-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a server process begins to write a dirty buffer. (If this happens often, it implies that shared_buffers is too small or the background writer control parameters need adjustment.) arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation.</td>
</tr>
<tr>
<td>buffer-write-dirty-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a dirty-buffer write is complete. The arguments are the same as for buffer-write-dirty-start.</td>
</tr>
<tr>
<td>wal-buffer-write-dirty-start</td>
<td>()</td>
<td>Probe that fires when a server process begins to write a dirty WAL buffer because no more WAL buffer space is available. (If this happens often, it implies that wal_buffers is too small.)</td>
</tr>
<tr>
<td>wal-buffer-write-dirty-done</td>
<td>()</td>
<td>Probe that fires when a dirty WAL buffer write is complete.</td>
</tr>
<tr>
<td>wal-insert</td>
<td>(unsigned char, unsigned char)</td>
<td>Probe that fires when a WAL record is inserted. arg0 is the resource manager (rmid) for the</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>------------------------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>wal-switch</td>
<td>()</td>
<td>Probe that fires when a WAL segment switch is requested.</td>
</tr>
<tr>
<td>smgr-md-read-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int)</td>
<td>Probe that fires when beginning to read a block from a relation. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer.</td>
</tr>
<tr>
<td>smgr-md-read-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, int, int)</td>
<td>Probe that fires when a block read is complete. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer. arg6 is the number of bytes actually read, while arg7 is the number requested (if these are different it indicates trouble).</td>
</tr>
<tr>
<td>smgr-md-write-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int)</td>
<td>Probe that fires when beginning to write a block to a relation. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer.</td>
</tr>
<tr>
<td>smgr-md-write-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, int, int)</td>
<td>Probe that fires when a block write is complete. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer.</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>-----------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>sort-start</td>
<td>(int, bool, int, int, bool, int)</td>
<td>Probe that fires when a sort operation is started. arg0 indicates heap, index or datum sort. arg1 is true for unique-value enforcement. arg2 is the number of key columns. arg3 is the number of kilobytes of work memory allowed. arg4 is true if random access to the sort result is required. arg5 indicates serial when 0, parallel worker when 1, or parallel leader when 2.</td>
</tr>
<tr>
<td>sort-done</td>
<td>(bool, long)</td>
<td>Probe that fires when a sort is complete. arg0 is true for external sort, false for internal sort. arg1 is the number of disk blocks used for an external sort, or kilobytes of memory used for an internal sort.</td>
</tr>
<tr>
<td>lwlock-acquire</td>
<td>(char *, LWLockMode)</td>
<td>Probe that fires when an LWLock has been acquired. arg0 is the LWLock's tranche. arg1 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-release</td>
<td>(char *)</td>
<td>Probe that fires when an LWLock has been released (but note that any released waiters have not yet been awakened). arg0 is the LWLock's tranche.</td>
</tr>
<tr>
<td>lwlock-wait-start</td>
<td>(char *, LWLockMode)</td>
<td>Probe that fires when an LWLock was not immediately available and a server process has begun to wait for the lock to become available. arg0 is the LWLock's tranche. arg1 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-wait-done</td>
<td>(char *, LWLockMode)</td>
<td>Probe that fires when a server process has been released from its wait for an LWLock (it does not actually have the lock yet). arg0 is the LWLock's tranche. arg1 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-condaquire</td>
<td>(char *, LWLockMode)</td>
<td>Probe that fires when an LWLock was successfully acquired when the caller specified no waiting. arg0 is the LWLock's tranche. arg1 is the requested lock mode, either exclusive or shared.</td>
</tr>
</tbody>
</table>
### Monitoring Database Activity

<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>lwlock-condacquire-fail</td>
<td>(char *, LWLockMode)</td>
<td>Probe that fires when an LWLock was not successfully acquired when the caller specified no waiting. arg0 is the LWLock's tranche. arg1 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lock-wait-start</td>
<td>(unsigned int, unsigned int, unsigned int, unsigned int, unsigned int, LOCKMODE)</td>
<td>Probe that fires when a request for a heavyweight lock (lmgr lock) has begun to wait because the lock is not available. arg0 through arg3 are the tag fields identifying the object being locked. arg4 indicates the type of object being locked. arg5 indicates the lock type being requested.</td>
</tr>
<tr>
<td>lock-wait-done</td>
<td>(unsigned int, unsigned int, unsigned int, unsigned int, unsigned int, LOCKMODE)</td>
<td>Probe that fires when a request for a heavyweight lock (lmgr lock) has finished waiting (i.e., has acquired the lock). The arguments are the same as for lock-wait-start.</td>
</tr>
<tr>
<td>deadlock-found</td>
<td>()</td>
<td>Probe that fires when a deadlock is found by the deadlock detector.</td>
</tr>
</tbody>
</table>

#### Table 28.24. Defined Types Used in Probe Parameters

<table>
<thead>
<tr>
<th>Type</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>LocalTransactionId</td>
<td>unsigned int</td>
</tr>
<tr>
<td>LWLockMode</td>
<td>int</td>
</tr>
<tr>
<td>LOCKMODE</td>
<td>int</td>
</tr>
<tr>
<td>BlockNumber</td>
<td>unsigned int</td>
</tr>
<tr>
<td>Oid</td>
<td>unsigned int</td>
</tr>
<tr>
<td>ForkNumber</td>
<td>int</td>
</tr>
<tr>
<td>bool</td>
<td>char</td>
</tr>
</tbody>
</table>

#### 28.5.3. Using Probes

The example below shows a DTrace script for analyzing transaction counts in the system, as an alternative to snapshotting `pg_stat_database` before and after a performance test:

```bash
#!/usr/sbin/dtrace -qs
postgresql$1:::transaction-start
{
    @start["Start"] = count();
    self->ts = timestamp;
}
```
Monitoring Database Activity

```d
postgresql$1:::transaction-abort
{
    @abort["Abort"] = count();
}

postgresql$1:::transaction-commit
/self->ts/
{
    @commit["Commit"] = count();
    @time["Total time (ns)"] = sum(timestamp - self->ts);
    self->ts=0;
}
```

When executed, the example D script gives output such as:

```
# ./txn_count.d `pgrep -n postgres` or ./txn_count.d <PID>
^C

Start 71
Commit 70
Total time (ns) 2312105013
```

Note

SystemTap uses a different notation for trace scripts than DTrace does, even though the underlying trace points are compatible. One point worth noting is that at this writing, SystemTap scripts must reference probe names using double underscores in place of hyphens. This is expected to be fixed in future SystemTap releases.

You should remember that DTrace scripts need to be carefully written and debugged, otherwise the trace information collected might be meaningless. In most cases where problems are found it is the instrumentation that is at fault, not the underlying system. When discussing information found using dynamic tracing, be sure to enclose the script used to allow that too to be checked and discussed.

### 28.5.4. Defining New Probes

New probes can be defined within the code wherever the developer desires, though this will require a recompilation. Below are the steps for inserting new probes:

1. Decide on probe names and data to be made available through the probes
2. Add the probe definitions to `src/backend/utils/probes.d`
3. Include `pg_trace.h` if it is not already present in the module(s) containing the probe points, and insert `TRACE_POSTGRESQL` probe macros at the desired locations in the source code
4. Recompile and verify that the new probes are available

**Example:** Here is an example of how you would add a probe to trace all new transactions by transaction ID.

1. Decide that the probe will be named `transaction-start` and requires a parameter of type `LocalTransactionId`
2. Add the probe definition to `src/backend/utils/probes.d`:
Monitoring Database Activity

probe transaction__start (LocalTransactionId);

Note the use of the double underline in the probe name. In a DTrace script using the probe, the double underline needs to be replaced with a hyphen, so transaction-start is the name to document for users.

3. At compile time, transaction__start is converted to a macro called TRACE_POSTGRESQL_TRANSACTION_START (notice the underscores are single here), which is available by including pg_trace.h. Add the macro call to the appropriate location in the source code. In this case, it looks like the following:

TRACE_POSTGRESQL_TRANSACTION_START(vxid.localTransactionId);

4. After recompiling and running the new binary, check that your newly added probe is available by executing the following DTrace command. You should see similar output:

    # dtrace -ln transaction-start
    ID   PROVIDER        MODULE           FUNCTION NAME
    18705 postgresql49878 postgres          StartTransactionCommand
        transaction-start
    18755 postgresql49877 postgres          StartTransactionCommand
        transaction-start
    18805 postgresql49876 postgres          StartTransactionCommand
        transaction-start
    18855 postgresql49875 postgres          StartTransactionCommand
        transaction-start
    18986 postgresql49873 postgres          StartTransactionCommand
        transaction-start

There are a few things to be careful about when adding trace macros to the C code:

• You should take care that the data types specified for a probe’s parameters match the data types of the variables used in the macro. Otherwise, you will get compilation errors.

• On most platforms, if PostgreSQL is built with --enable-dtrace, the arguments to a trace macro will be evaluated whenever control passes through the macro, even if no tracing is being done. This is usually not worth worrying about if you are just reporting the values of a few local variables. But beware of putting expensive function calls into the arguments. If you need to do that, consider protecting the macro with a check to see if the trace is actually enabled:

    if (TRACE_POSTGRESQL_TRANSACTION_START_ENABLED())
        TRACE_POSTGRESQL_TRANSACTION_START(some_function(...));

Each trace macro has a corresponding ENABLED macro.
Chapter 29. Monitoring Disk Usage

This chapter discusses how to monitor the disk usage of a PostgreSQL database system.

29.1. Determining Disk Usage

Each table has a primary heap disk file where most of the data is stored. If the table has any columns with potentially-wide values, there also might be a TOAST file associated with the table, which is used to store values too wide to fit comfortably in the main table (see Section 68.2). There will be one valid index on the TOAST table, if present. There also might be indexes associated with the base table. Each table and index is stored in a separate disk file — possibly more than one file, if the file would exceed one gigabyte. Naming conventions for these files are described in Section 68.1.

You can monitor disk space in three ways: using the SQL functions listed in Table 9.84, using the oid2name module, or using manual inspection of the system catalogs. The SQL functions are the easiest to use and are generally recommended. The remainder of this section shows how to do it by inspection of the system catalogs.

Using psql on a recently vacuumed or analyzed database, you can issue queries to see the disk usage of any table:

```sql
SELECT pg_relation_filepath(oid), relpages FROM pg_class WHERE relname = 'customer';
```

<table>
<thead>
<tr>
<th>pg_relation_filepath</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>base/16384/16806</td>
<td>60</td>
</tr>
</tbody>
</table>

Each page is typically 8 kilobytes. (Remember, relpages is only updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX.) The file path name is of interest if you want to examine the table's disk file directly.

To show the space used by TOAST tables, use a query like the following:

```sql
SELECT relname, relpages FROM pg_class,
  (SELECT reltoastrelid
   FROM pg_class
   WHERE relname = 'customer') AS ss
WHERE oid = ss.reltoastrelid OR
  oid = (SELECT indexrelid
         FROM pg_index
         WHERE indrelid = ss.reltoastrelid)
ORDER BY relname;
```

<table>
<thead>
<tr>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_toast_16806</td>
<td>0</td>
</tr>
<tr>
<td>pg_toast_16806_index</td>
<td>1</td>
</tr>
</tbody>
</table>

You can easily display index sizes, too:

```sql
SELECT c2.relname, c2.relpages
```

738
FROM pg_class c, pg_class c2, pg_index i  
WHERE c.relname = 'customer' AND  
    c.oid = i.indrelid AND  
    c2.oid = i.indexrelid  
ORDER BY c2.relname;

<table>
<thead>
<tr>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>customer_id_indexdex</td>
<td>26</td>
</tr>
</tbody>
</table>

It is easy to find your largest tables and indexes using this information:

```
SELECT relname, relpages  
FROM pg_class  
ORDER BY relpages DESC;
```

<table>
<thead>
<tr>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>bigtable</td>
<td>3290</td>
</tr>
<tr>
<td>customer</td>
<td>3144</td>
</tr>
</tbody>
</table>

### 29.2. Disk Full Failure

The most important disk monitoring task of a database administrator is to make sure the disk doesn’t become full. A filled data disk will not result in data corruption, but it might prevent useful activity from occurring. If the disk holding the WAL files grows full, database server panic and consequent shutdown might occur.

If you cannot free up additional space on the disk by deleting other things, you can move some of the database files to other file systems by making use of tablespaces. See Section 22.6 for more information about that.

**Tip**

Some file systems perform badly when they are almost full, so do not wait until the disk is completely full to take action.

If your system supports per-user disk quotas, then the database will naturally be subject to whatever quota is placed on the user the server runs as. Exceeding the quota will have the same bad effects as running out of disk space entirely.
Chapter 30. Reliability and the Write-Ahead Log

This chapter explains how the Write-Ahead Log is used to obtain efficient, reliable operation.

30.1. Reliability

Reliability is an important property of any serious database system, and PostgreSQL does everything possible to guarantee reliable operation. One aspect of reliable operation is that all data recorded by a committed transaction should be stored in a nonvolatile area that is safe from power loss, operating system failure, and hardware failure (except failure of the nonvolatile area itself, of course). Successfully writing the data to the computer's permanent storage (disk drive or equivalent) ordinarily meets this requirement. In fact, even if a computer is fatally damaged, if the disk drives survive they can be moved to another computer with similar hardware and all committed transactions will remain intact.

While forcing data to the disk platters periodically might seem like a simple operation, it is not. Because disk drives are dramatically slower than main memory and CPUs, several layers of caching exist between the computer's main memory and the disk platters. First, there is the operating system's buffer cache, which caches frequently requested disk blocks and combines disk writes. Fortunately, all operating systems give applications a way to force writes from the buffer cache to disk, and PostgreSQL uses those features. (See the wal_sync_method parameter to adjust how this is done.)

Next, there might be a cache in the disk drive controller; this is particularly common on RAID controller cards. Some of these caches are write-through, meaning writes are sent to the drive as soon as they arrive. Others are write-back, meaning data is sent to the drive at some later time. Such caches can be a reliability hazard because the memory in the disk controller cache is volatile, and will lose its contents in a power failure. Better controller cards have battery-backup units (BBUs), meaning the card has a battery that maintains power to the cache in case of system power loss. After power is restored the data will be written to the disk drives.

And finally, most disk drives have caches. Some are write-through while some are write-back, and the same concerns about data loss exist for write-back drive caches as for disk controller caches. Consumer-grade IDE and SATA drives are particularly likely to have write-back caches that will not survive a power failure. Many solid-state drives (SSD) also have volatile write-back caches.

These caches can typically be disabled; however, the method for doing this varies by operating system and drive type:

- On Linux, IDE and SATA drives can be queried using hdparm -I; write caching is enabled if there is a * next to write cache. hdparm -W 0 can be used to turn off write caching. SCSI drives can be queried using sdparm. Use sdparm --get=WCE to check whether the write cache is enabled and sdparm --clear=WCE to disable it.

- On FreeBSD, IDE drives can be queried using atacontrol and write caching turned off using hw.ata.wc=0 in /boot/loader.conf; SCSI drives can be queried using camcontrol identify, and the write cache both queried and changed using sdparm when available.

- On Solaris, the disk write cache is controlled by format -e. (The Solaris ZFS file system is safe with disk write-cache enabled because it issues its own disk cache flush commands.)

- On Windows, if wal_sync_method is open_datsasync (the default), write caching can be disabled by unchecking My Computer\Open\disk drive\Properties\Hardware\Properties\Policies\Enable write caching on the disk. Alternatively, set wal_sync_method to fsync or fsync_writethrough, which prevent write caching.

1 http://sg.danny.cz/sg/sdparm.html
Reliability and the Write-Ahead Log

- On macOS, write caching can be prevented by setting `wal_sync_method` to `fsync_writethrough`.

Recent SATA drives (those following ATAPI-6 or later) offer a drive cache flush command (FLUSH CACHE EXT), while SCSI drives have long supported a similar command SYNCHRONIZE CACHE. These commands are not directly accessible to PostgreSQL, but some file systems (e.g., ZFS, ext4) can use them to flush data to the platters on write-back-enabled drives. Unfortunately, such file systems behave suboptimally when combined with battery-backup unit (BBU) disk controllers. In such setups, the synchronize command forces all data from the controller cache to the disks, eliminating much of the benefit of the BBU. You can run the `pg_test_fsync` program to see if you are affected. If you are affected, the performance benefits of the BBU can be regained by turning off write barriers in the file system or reconfiguring the disk controller, if that is an option. If write barriers are turned off, make sure the battery remains functional; a faulty battery can potentially lead to data loss. Hopefully file system and disk controller designers will eventually address this suboptimal behavior.

When the operating system sends a write request to the storage hardware, there is little it can do to make sure the data has arrived at a truly non-volatile storage area. Rather, it is the administrator's responsibility to make certain that all storage components ensure integrity for both data and file-system metadata. Avoid disk controllers that have non-battery-backed write caches. At the drive level, disable write-back caching if the drive cannot guarantee the data will be written before shutdown. If you use SSDs, be aware that many of these do not honor cache flush commands by default. You can test for reliable I/O subsystem behavior using `diskchecker.pl`.

Another risk of data loss is posed by the disk platter write operations themselves. Disk platters are divided into sectors, commonly 512 bytes each. Every physical read or write operation processes a whole sector. When a write request arrives at the drive, it might be for some multiple of 512 bytes (PostgreSQL typically writes 8192 bytes, or 16 sectors, at a time), and the process of writing could fail due to power loss at any time, meaning some of the 512-byte sectors were written while others were not. To guard against such failures, PostgreSQL periodically writes full page images to permanent WAL storage before modifying the actual page on disk. By doing this, during crash recovery PostgreSQL can restore partially-written pages from WAL. If you have file-system software that prevents partial page writes (e.g., ZFS), you can turn off this page imaging by turning off the `full_page_writes` parameter. Battery-Backed Unit (BBU) disk controllers do not prevent partial page writes unless they guarantee that data is written to the BBU as full (8kB) pages.

PostgreSQL also protects against some kinds of data corruption on storage devices that may occur because of hardware errors or media failure over time, such as reading/writing garbage data.

- Each individual record in a WAL file is protected by a CRC-32 (32-bit) check that allows us to tell if record contents are correct. The CRC value is set when we write each WAL record and checked during crash recovery, archive recovery and replication.

- Data pages are not currently checksummed by default, though full page images recorded in WAL records will be protected; see `initdb` for details about enabling data page checksums.

- Internal data structures such as `pg_xact`, `pg_subtrans`, `pg_multixact`, `pg_serial`, `pg_notify`, `pg_stat`, `pg_snapshots` are not directly checksummed, nor are pages protected by full page writes. However, where such data structures are persistent, WAL records are written that allow recent changes to be accurately rebuilt at crash recovery and those WAL records are protected as discussed above.

- Individual state files in `pg_twophase` are protected by CRC-32.

- Temporary data files used in larger SQL queries for sorts, materializations and intermediate results are not currently checksummed, nor will WAL records be written for changes to those files.

PostgreSQL does not protect against correctable memory errors and it is assumed you will operate using RAM that uses industry standard Error Correcting Codes (ECC) or better protection.

---

2 https://brad.livejournal.com/2116715.html
30.2. Write-Ahead Logging (WAL)

*Write-Ahead Logging* (WAL) is a standard method for ensuring data integrity. A detailed description can be found in most (if not all) books about transaction processing. Briefly, WAL’s central concept is that changes to data files (where tables and indexes reside) must be written only after those changes have been logged, that is, after log records describing the changes have been flushed to permanent storage. If we follow this procedure, we do not need to flush data pages to disk on every transaction commit, because we know that in the event of a crash we will be able to recover the database using the log: any changes that have not been applied to the data pages can be redone from the log records. (This is roll-forward recovery, also known as REDO.)

**Tip**

Because WAL restores database file contents after a crash, journaled file systems are not necessary for reliable storage of the data files or WAL files. In fact, journaling overhead can reduce performance, especially if journaling causes file system data to be flushed to disk. Fortunately, data flushing during journaling can often be disabled with a file system mount option, e.g. `data=writeback` on a Linux ext3 file system. Journaled file systems do improve boot speed after a crash.

Using WAL results in a significantly reduced number of disk writes, because only the log file needs to be flushed to disk to guarantee that a transaction is committed, rather than every data file changed by the transaction. The log file is written sequentially, and so the cost of syncing the log is much less than the cost of flushing the data pages. This is especially true for servers handling many small transactions touching different parts of the data store. Furthermore, when the server is processing many small concurrent transactions, one `fsync` of the log file may suffice to commit many transactions.

WAL also makes it possible to support on-line backup and point-in-time recovery, as described in Section 25.3. By archiving the WAL data we can support reverting to any time instant covered by the available WAL data: we simply install a prior physical backup of the database, and replay the WAL log just as far as the desired time. What's more, the physical backup doesn't have to be an instantaneous snapshot of the database state — if it is made over some period of time, then replaying the WAL log for that period will fix any internal inconsistencies.

30.3. Asynchronous Commit

*Asynchronous commit* is an option that allows transactions to complete more quickly, at the cost that the most recent transactions may be lost if the database should crash. In many applications this is an acceptable trade-off.

As described in the previous section, transaction commit is normally *synchronous*: the server waits for the transaction's WAL records to be flushed to permanent storage before returning a success indication to the client. The client is therefore guaranteed that a transaction reported to be committed will be preserved, even in the event of a server crash immediately after. However, for short transactions this delay is a major component of the total transaction time. Selecting asynchronous commit mode means that the server returns success as soon as the transaction is logically completed, before the WAL records it generated have actually made their way to disk. This can provide a significant boost in throughput for small transactions.

Asynchronous commit introduces the risk of data loss. There is a short time window between the report of transaction completion to the client and the time that the transaction is truly committed (that is, it is guaranteed not to be lost if the server crashes). Thus asynchronous commit should not be used if the client will take external actions relying on the assumption that the transaction will be remembered. As an example, a bank would certainly not use asynchronous commit for a transaction recording an ATM’s dispensing of cash. But in many scenarios, such as event logging, there is no need for a strong guarantee of this kind.
The risk that is taken by using asynchronous commit is of data loss, not data corruption. If the database should crash, it will recover by replaying WAL up to the last record that was flushed. The database will therefore be restored to a self-consistent state, but any transactions that were not yet flushed to disk will not be reflected in that state. The net effect is therefore loss of the last few transactions. Because the transactions are replayed in commit order, no inconsistency can be introduced — for example, if transaction B made changes relying on the effects of a previous transaction A, it is not possible for A's effects to be lost while B's effects are preserved.

The user can select the commit mode of each transaction, so that it is possible to have both synchronous and asynchronous commit transactions running concurrently. This allows flexible trade-offs between performance and certainty of transaction durability. The commit mode is controlled by the user-settable parameter synchronous_commit, which can be changed in any of the ways that a configuration parameter can be set. The mode used for any one transaction depends on the value of synchronous_commit when transaction commit begins.

Certain utility commands, for instance DROP TABLE, are forced to commit synchronously regardless of the setting of synchronous_commit. This is to ensure consistency between the server's file system and the logical state of the database. The commands supporting two-phase commit, such as PREPARE TRANSACTION, are also always synchronous.

If the database crashes during the risk window between an asynchronous commit and the writing of the transaction's WAL records, then changes made during that transaction will be lost. The duration of the risk window is limited because a background process (the “WAL writer”) flushes unwritten WAL records to disk every wal_writer_delay milliseconds. The actual maximum duration of the risk window is three times wal_writer_delay because the WAL writer is designed to favor writing whole pages at a time during busy periods.

**Caution**

An immediate-mode shutdown is equivalent to a server crash, and will therefore cause loss of any unflushed asynchronous commits.

Asynchronous commit provides behavior different from setting fsync = off. fsync is a server-wide setting that will alter the behavior of all transactions. It disables all logic within PostgreSQL that attempts to synchronize writes to different portions of the database, and therefore a system crash (that is, a hardware or operating system crash, not a failure of PostgreSQL itself) could result in arbitrarily bad corruption of the database state. In many scenarios, asynchronous commit provides most of the performance improvement that could be obtained by turning off fsync, but without the risk of data corruption.

commit_delay also sounds very similar to asynchronous commit, but it is actually a synchronous commit method (in fact, commit_delay is ignored during an asynchronous commit). commit_delay causes a delay just before a transaction flushes WAL to disk, in the hope that a single flush executed by one such transaction can also serve other transactions committing at about the same time. The setting can be thought of as a way of increasing the time window in which transactions can join a group about to participate in a single flush, to amortize the cost of the flush among multiple transactions.

### 30.4. WAL Configuration

There are several WAL-related configuration parameters that affect database performance. This section explains their use. Consult Chapter 19 for general information about setting server configuration parameters.

*Checkpoints* are points in the sequence of transactions at which it is guaranteed that the heap and index data files have been updated with all information written before that checkpoint. At checkpoint time, all dirty data pages are flushed to disk and a special checkpoint record is written to the log file. (The
Reliability and the Write-Ahead Log

change records were previously flushed to the WAL files.) In the event of a crash, the crash recovery procedure looks at the latest checkpoint record to determine the point in the log (known as the redo record) from which it should start the REDO operation. Any changes made to data files before that point are guaranteed to be already on disk. Hence, after a checkpoint, log segments preceding the one containing the redo record are no longer needed and can be recycled or removed. (When WAL archiving is being done, the log segments must be archived before being recycled or removed.)

The checkpoint requirement of flushing all dirty data pages to disk can cause a significant I/O load. For this reason, checkpoint activity is throttled so that I/O begins at checkpoint start and completes before the next checkpoint is due to start; this minimizes performance degradation during checkpoints.

The server's checkpointer process automatically performs a checkpoint every so often. A checkpoint is begun every checkpoint_timeout seconds, or if max_wal_size is about to be exceeded, whichever comes first. The default settings are 5 minutes and 1 GB, respectively. If no WAL has been written since the previous checkpoint, new checkpoints will be skipped even if checkpoint_timeout has passed. (If WAL archiving is being used and you want to put a lower limit on how often files are archived in order to bound potential data loss, you should adjust the archive_timeout parameter rather than the checkpoint parameters.) It is also possible to force a checkpoint by using the SQL command CHECKPOINT.

Reducing checkpoint_timeout and/or max_wal_size causes checkpoints to occur more often. This allows faster after-crash recovery, since less work will need to be redone. However, one must balance this against the increased cost of flushing dirty data pages more often. If full_page_writes is set (as is the default), there is another factor to consider. To ensure data page consistency, the first modification of a data page after each checkpoint results in logging the entire page content. In that case, a smaller checkpoint interval increases the volume of output to the WAL log, partially negating the goal of using a smaller interval, and in any case causing more disk I/O.

Checkpoints are fairly expensive, first because they require writing out all currently dirty buffers, and second because they result in extra subsequent WAL traffic as discussed above. It is therefore wise to set the checkpointing parameters high enough so that checkpoints don't happen too often. As a simple sanity check on your checkpointing parameters, you can set the checkpoint_warning parameter. If checkpoints happen closer together than checkpoint_warning seconds, a message will be output to the server log recommending increasing max_wal_size. Occasional appearance of such a message is not cause for alarm, but if it appears often then the checkpoint control parameters should be increased. Bulk operations such as large COPY transfers might cause a number of such warnings to appear if you have not set max_wal_size high enough.

To avoid flooding the I/O system with a burst of page writes, writing dirty buffers during a checkpoint is spread over a period of time. That period is controlled by checkpoint_completion_target, which is given as a fraction of the checkpoint interval. The I/O rate is adjusted so that the checkpoint finishes when the given fraction of checkpoint_timeout seconds have elapsed, or before max_wal_size is exceeded, whichever is sooner. With the default value of 0.5, PostgreSQL can be expected to complete each checkpoint in about half the time before the next checkpoint starts. On a system that's very close to maximum I/O throughput during normal operation, you might want to increase checkpoint_completion_target to reduce the I/O load from checkpoints. The disadvantage of this is that prolonging checkpoints affects recovery time, because more WAL segments will need to be kept around for possible use in recovery. Although checkpoint_completion_target can be set as high as 1.0, it is best to keep it less than that (perhaps 0.9 at most) since checkpoints include some other activities besides writing dirty buffers. A setting of 1.0 is quite likely to result in checkpoints not being completed on time, which would result in performance loss due to unexpected variation in the number of WAL segments needed.

On Linux and POSIX platforms checkpoint_flush_after allows to force the OS that pages written by the checkpoint should be flushed to disk after a configurable number of bytes. Otherwise, these pages may be kept in the OS's page cache, inducing a stall when fsync is issued at the end of a checkpoint. This setting will often help to reduce transaction latency, but it also can have an adverse effect on performance; particularly for workloads that are bigger than shared_buffers, but smaller than the OS's page cache.
The number of WAL segment files in the `pg_wal` directory depends on `min_wal_size`, `max_wal_size`, and the amount of WAL generated in previous checkpoint cycles. When old log segment files are no longer needed, they are removed or recycled (that is, renamed to become future segments in the numbered sequence). If, due to a short-term peak of log output rate, `max_wal_size` is exceeded, the unneeded segment files will be removed until the system gets back under this limit. Below that limit, the system recycles enough WAL files to cover the estimated need until the next checkpoint, and removes the rest. The estimate is based on a moving average of the number of WAL files used in previous checkpoint cycles. The moving average is increased immediately if the actual usage exceeds the estimate, so it accommodates peak usage rather than average usage to some extent. `min_wal_size` puts a minimum on the amount of WAL files recycled for future usage; that much WAL is always recycled for future use, even if the system is idle and the WAL usage estimate suggests that little WAL is needed.

Independently of `max_wal_size`, `wal_keep_segments + 1` most recent WAL files are kept at all times. Also, if WAL archiving is used, old segments can not be removed or recycled until they are archived. If WAL archiving cannot keep up with the pace that WAL is generated, or if `archive_command` fails repeatedly, old WAL files will accumulate in `pg_wal` until the situation is resolved. A slow or failed standby server that uses a replication slot will have the same effect (see Section 26.2.6).

In archive recovery or standby mode, the server periodically performs restartpoints, which are similar to checkpoints in normal operation: the server forces all its state to disk, updates the `pg_control` file to indicate that the already-processed WAL data need not be scanned again, and then recycles any old log segment files in the `pg_wal` directory. Restartpoints can't be performed more frequently than checkpoints in the master because restartpoints can only be performed at checkpoint records. A restartpoint is triggered when a checkpoint record is reached if at least `checkpoint_timeout` seconds have passed since the last restartpoint, or if WAL size is about to exceed `max_wal_size`. However, because of limitations on when a restartpoint can be performed, `max_wal_size` is often exceeded during recovery, by up to one checkpoint cycle's worth of WAL. (`max_wal_size` is never a hard limit anyway, so you should always leave plenty of headroom to avoid running out of disk space.)

There are two commonly used internal WAL functions: `XLogInsertRecord` and `XLogFlush`. `XLogInsertRecord` is used to place a new record into the WAL buffers in shared memory. If there is no space for the new record, `XLogInsertRecord` will have to write (move to kernel cache) a few filled WAL buffers. This is undesirable because `XLogInsertRecord` is used on every database low level modification (for example, row insertion) at a time when an exclusive lock is held on affected data pages, so the operation needs to be as fast as possible. What is worse, writing WAL buffers might also force the creation of a new log segment, which takes even more time. Normally, WAL buffers should be written and flushed by an `XLogFlush` request, which is made, for the most part, at transaction commit time to ensure that transaction records are flushed to permanent storage. On systems with high log output, `XLogFlush` requests might not occur often enough to prevent `XLogInsertRecord` from having to do writes. On such systems one should increase the number of WAL buffers by modifying the `wal_buffers` parameter. When `full_page_writes` is set and the system is very busy, setting `wal_buffers` higher will help smooth response times during the period immediately following each checkpoint.

The `commit_delay` parameter defines for how many microseconds a group commit leader process will sleep after acquiring a lock within `XLogFlush`, while group commit followers queue up behind the leader. This delay allows other server processes to add their commit records to the WAL buffers so that all of them will be flushed by the leader's eventual sync operation. No sleep will occur if fsync is not enabled, or if fewer than `commit_siblings` other sessions are currently in active transactions; this avoids sleeping when it's unlikely that any other session will commit soon. Note that on some platforms, the resolution of a sleep request is ten milliseconds, so that any nonzero `commit_delay` setting between 1 and 10000 microseconds would have the same effect. Note also that on some platforms, sleep operations may take slightly longer than requested by the parameter.

Since the purpose of `commit_delay` is to allow the cost of each flush operation to be amortized across concurrently committing transactions (potentially at the expense of transaction latency), it is
Reliability and the Write-Ahead Log

necessary to quantify that cost before the setting can be chosen intelligently. The higher that cost is, the more effective commit_delay is expected to be in increasing transaction throughput, up to a point. The pg_test_fsync program can be used to measure the average time in microseconds that a single WAL flush operation takes. A value of half of the average time the program reports it takes to flush after a single 8kB write operation is often the most effective setting for commit_delay, so this value is recommended as the starting point to use when optimizing for a particular workload. While tuning commit_delay is particularly useful when the WAL log is stored on high-latency rotating disks, benefits can be significant even on storage media with very fast sync times, such as solid-state drives or RAID arrays with a battery-backed write cache; but this should definitely be tested against a representative workload. Higher values of commit_siblings should be used in such cases, whereas smaller commit_siblings values are often helpful on higher latency media. Note that it is quite possible that a setting of commit_delay that is too high can increase transaction latency by so much that total transaction throughput suffers.

When commit_delay is set to zero (the default), it is still possible for a form of group commit to occur, but each group will consist only of sessions that reach the point where they need to flush their commit records during the window in which the previous flush operation (if any) is occurring. At higher client counts a “gangway effect” tends to occur, so that the effects of group commit become significant even when commit_delay is zero, and thus explicitly setting commit_delay tends to help less. Setting commit_delay can only help when (1) there are some concurrently committing transactions, and (2) throughput is limited to some degree by commit rate; but with high rotational latency this setting can be effective in increasing transaction throughput with as few as two clients (that is, a single committing client with one sibling transaction).

The wal_sync_method parameter determines how PostgreSQL will ask the kernel to force WAL updates out to disk. All the options should be the same in terms of reliability, with the exception of fsync_writethrough, which can sometimes force a flush of the disk cache even when other options do not do so. However, it's quite platform-specific which one will be the fastest. You can test the speeds of different options using the pg_test_fsync program. Note that this parameter is irrelevant if fsync has been turned off.

Enabling the wal_debug configuration parameter (provided that PostgreSQL has been compiled with support for it) will result in each XLogInsertRecord and XLogFlush WAL call being logged to the server log. This option might be replaced by a more general mechanism in the future.

30.5. WAL Internals

WAL is automatically enabled; no action is required from the administrator except ensuring that the disk-space requirements for the WAL logs are met, and that any necessary tuning is done (see Section 30.4).

WAL records are appended to the WAL logs as each new record is written. The insert position is described by a Log Sequence Number (LSN) that is a byte offset into the logs, increasing monotonically with each new record. LSN values are returned as the datatype pg_lsn. Values can be compared to calculate the volume of WAL data that separates them, so they are used to measure the progress of replication and recovery.

WAL logs are stored in the directory pg_wal under the data directory, as a set of segment files, normally each 16 MB in size (but the size can be changed by altering the --wal-segsize initdb option). Each segment is divided into pages, normally 8 kB each (this size can be changed via the --with-wal-blocksize configure option). The log record headers are described in access/xlocrecord.h; the record content is dependent on the type of event that is being logged. Segment files are given ever-increasing numbers as names, starting at 00000001000000000000000000000000. The numbers do not wrap, but it will take a very, very long time to exhaust the available stock of numbers.

It is advantageous if the log is located on a different disk from the main database files. This can be achieved by moving the pg_wal directory to another location (while the server is shut down, of course) and creating a symbolic link from the original location in the main data directory to the new location.
Reliability and the Write-Ahead Log

The aim of WAL is to ensure that the log is written before database records are altered, but this can be subverted by disk drives that falsely report a successful write to the kernel, when in fact they have only cached the data and not yet stored it on the disk. A power failure in such a situation might lead to irrecoverable data corruption. Administrators should try to ensure that disks holding PostgreSQL's WAL log files do not make such false reports. (See Section 30.1.)

After a checkpoint has been made and the log flushed, the checkpoint's position is saved in the file pg_control. Therefore, at the start of recovery, the server first reads pg_control and then the checkpoint record; then it performs the REDO operation by scanning forward from the log location indicated in the checkpoint record. Because the entire content of data pages is saved in the log on the first page modification after a checkpoint (assuming full_page_writes is not disabled), all pages changed since the checkpoint will be restored to a consistent state.

To deal with the case where pg_control is corrupt, we should support the possibility of scanning existing log segments in reverse order — newest to oldest — in order to find the latest checkpoint. This has not been implemented yet. pg_control is small enough (less than one disk page) that it is not subject to partial-write problems, and as of this writing there have been no reports of database failures due solely to the inability to read pg_control itself. So while it is theoretically a weak spot, pg_control does not seem to be a problem in practice.
Chapter 31. Logical Replication

Logical replication is a method of replicating data objects and their changes, based upon their replication identity (usually a primary key). We use the term logical in contrast to physical replication, which uses exact block addresses and byte-by-byte replication. PostgreSQL supports both mechanisms concurrently, see Chapter 26. Logical replication allows fine-grained control over both data replication and security.

Logical replication uses a publish and subscribe model with one or more subscribers subscribing to one or more publications on a publisher node. Subscribers pull data from the publications they subscribe to and may subsequently re-publish data to allow cascading replication or more complex configurations.

Logical replication of a table typically starts with taking a snapshot of the data on the publisher database and copying that to the subscriber. Once that is done, the changes on the publisher are sent to the subscriber as they occur in real-time. The subscriber applies the data in the same order as the publisher so that transactional consistency is guaranteed for publications within a single subscription. This method of data replication is sometimes referred to as transactional replication.

The typical use-cases for logical replication are:

• Sending incremental changes in a single database or a subset of a database to subscribers as they occur.
• Firing triggers for individual changes as they arrive on the subscriber.
• Consolidating multiple databases into a single one (for example for analytical purposes).
• Replicating between different major versions of PostgreSQL.
• Replicating between PostgreSQL instances on different platforms (for example Linux to Windows)
• Giving access to replicated data to different groups of users.
• Sharing a subset of the database between multiple databases.

The subscriber database behaves in the same way as any other PostgreSQL instance and can be used as a publisher for other databases by defining its own publications. When the subscriber is treated as read-only by application, there will be no conflicts from a single subscription. On the other hand, if there are other writes done either by an application or by other subscribers to the same set of tables, conflicts can arise.

31.1. Publication

A publication can be defined on any physical replication master. The node where a publication is defined is referred to as publisher. A publication is a set of changes generated from a table or a group of tables, and might also be described as a change set or replication set. Each publication exists in only one database.

Publications are different from schemas and do not affect how the table is accessed. Each table can be added to multiple publications if needed. Publications may currently only contain tables. Objects must be added explicitly, except when a publication is created for ALL TABLES.

Publications can choose to limit the changes they produce to any combination of INSERT, UPDATE, DELETE, and TRUNCATE, similar to how triggers are fired by particular event types. By default, all operation types are replicated.

A published table must have a “replica identity” configured in order to be able to replicate UPDATE and DELETE operations, so that appropriate rows to update or delete can be identified on the subscriber side. By default, this is the primary key, if there is one. Another unique index (with certain additional requirements) can also be set to be the replica identity. If the table does not have any suitable key, then
it can be set to replica identity "full", which means the entire row becomes the key. This, however,
is very inefficient and should only be used as a fallback if no other solution is possible. If a replica
identity other than "full" is set on the publisher side, a replica identity comprising the same or fewer
columns must also be set on the subscriber side. See REPLICA IDENTITY for details on how to
set the replica identity. If a table without a replica identity is added to a publication that replicates
UPDATE or DELETE operations then subsequent UPDATE or DELETE operations will cause an error
on the publisher. INSERT operations can proceed regardless of any replica identity.

Every publication can have multiple subscribers.

A publication is created using the CREATE PUBLICATION command and may later be altered or
dropped using corresponding commands.

The individual tables can be added and removed dynamically using ALTER PUBLICATION. Both the
ADD TABLE and DROP TABLE operations are transactional; so the table will start or stop replicating
at the correct snapshot once the transaction has committed.

### 31.2. Subscription

A subscription is the downstream side of logical replication. The node where a subscription is defined
is referred to as the subscriber. A subscription defines the connection to another database and set of
publications (one or more) to which it wants to subscribe.

The subscriber database behaves in the same way as any other PostgreSQL instance and can be used
as a publisher for other databases by defining its own publications.

A subscriber node may have multiple subscriptions if desired. It is possible to define multiple sub-
scriptions between a single publisher-subscriber pair, in which case care must be taken to ensure that
the subscribed publication objects don't overlap.

Each subscription will receive changes via one replication slot (see Section 26.2.6). Additional tem-
porary replication slots may be required for the initial data synchronization of pre-existing table data.

A logical replication subscription can be a standby for synchronous replication (see Section 26.2.8).
The standby name is by default the subscription name. An alternative name can be specified as ap-
plication_name in the connection information of the subscription.

Subscriptions are dumped by pg_dump if the current user is a superuser. Otherwise a warning is
written and subscriptions are skipped, because non-superusers cannot read all subscription information
from the pg_subscription catalog.

The subscription is added using CREATE SUBSCRIPTION and can be stopped/resumed at any time
using the ALTER SUBSCRIPTION command and removed using DROP SUBSCRIPTION.

When a subscription is dropped and recreated, the synchronization information is lost. This means
that the data has to be resynchronized afterwards.

The schema definitions are not replicated, and the published tables must exist on the subscriber. Only
regular tables may be the target of replication. For example, you can't replicate to a view.

The tables are matched between the publisher and the subscriber using the fully qualified table name.
Replication to differently-named tables on the subscriber is not supported.

Columns of a table are also matched by name. A different order of columns in the target table is al-
lowed, but the column types have to match. The target table can have additional columns not provided
by the published table. Those will be filled with their default values.

### 31.2.1. Replication Slot Management

As mentioned earlier, each (active) subscription receives changes from a replication slot on the remote
(publishing) side. Normally, the remote replication slot is created automatically when the subscription
Logical Replication

is created using `CREATE SUBSCRIPTION` and it is dropped automatically when the subscription is dropped using `DROP SUBSCRIPTION`. In some situations, however, it can be useful or necessary to manipulate the subscription and the underlying replication slot separately. Here are some scenarios:

- When creating a subscription, the replication slot already exists. In that case, the subscription can be created using the `create_slot = false` option to associate with the existing slot.

- When creating a subscription, the remote host is not reachable or in an unclear state. In that case, the subscription can be created using the `connect = false` option. The remote host will then not be contacted at all. This is what `pg_dump` uses. The remote replication slot will then have to be created manually before the subscription can be activated.

- When dropping a subscription, the replication slot should be kept. This could be useful when the subscriber database is being moved to a different host and will be activated from there. In that case, disassociate the slot from the subscription using `ALTER SUBSCRIPTION` before attempting to drop the subscription.

- When dropping a subscription, the remote host is not reachable. In that case, disassociate the slot from the subscription using `ALTER SUBSCRIPTION` before attempting to drop the subscription. If the remote database instance no longer exists, no further action is then necessary. If, however, the remote database instance is just unreachable, the replication slot should then be dropped manually; otherwise it would continue to reserve WAL and might eventually cause the disk to fill up. Such cases should be carefully investigated.

31.3. Conflicts

Logical replication behaves similarly to normal DML operations in that the data will be updated even if it was changed locally on the subscriber node. If incoming data violates any constraints the replication will stop. This is referred to as a *conflict*. When replicating `UPDATE` or `DELETE` operations, missing data will not produce a conflict and such operations will simply be skipped.

A conflict will produce an error and will stop the replication; it must be resolved manually by the user. Details about the conflict can be found in the subscriber's server log.

The resolution can be done either by changing data on the subscriber so that it does not conflict with the incoming change or by skipping the transaction that conflicts with the existing data. The transaction can be skipped by calling the `pg_replication_origin_advance()` function with a `node_name` corresponding to the subscription name, and a position. The current position of origins can be seen in the `pg_replication_origin_status` system view.

31.4. Restrictions

Logical replication currently has the following restrictions or missing functionality. These might be addressed in future releases.

- The database schema and DDL commands are not replicated. The initial schema can be copied by hand using `pg_dump --schema-only`. Subsequent schema changes would need to be kept in sync manually. (Note, however, that there is no need for the schemas to be absolutely the same on both sides.) Logical replication is robust when schema definitions change in a live database: When the schema is changed on the publisher and replicated data starts arriving at the subscriber but does not fit into the table schema, replication will error until the schema is updated. In many cases, intermittent errors can be avoided by applying additive schema changes to the subscriber first.

- Sequence data is not replicated. The data in serial or identity columns backed by sequences will of course be replicated as part of the table, but the sequence itself would still show the start value on the subscriber. If the subscriber is used as a read-only database, then this should typically not be a problem. If, however, some kind of switchover or failover to the subscriber database is intended, then the sequences would need to be updated to the latest values, either by copying the current data
Logical Replication

from the publisher (perhaps using pg_dump) or by determining a sufficiently high value from the tables themselves.

• Replication of TRUNCATE commands is supported, but some care must be taken when truncating groups of tables connected by foreign keys. When replicating a truncate action, the subscriber will truncate the same group of tables that was truncated on the publisher, either explicitly specified or implicitly collected via CASCADE, minus tables that are not part of the subscription. This will work correctly if all affected tables are part of the same subscription. But if some tables to be truncated on the subscriber have foreign-key links to tables that are not part of the same (or any) subscription, then the application of the truncate action on the subscriber will fail.

• Large objects (see Chapter 35) are not replicated. There is no workaround for that, other than storing data in normal tables.

• Replication is only possible from base tables to base tables. That is, the tables on the publication and on the subscription side must be normal tables, not views, materialized views, partition root tables, or foreign tables. In the case of partitions, you can therefore replicate a partition hierarchy one-to-one, but you cannot currently replicate to a differently partitioned setup. Attempts to replicate tables other than base tables will result in an error.

31.5. Architecture

Logical replication starts by copying a snapshot of the data on the publisher database. Once that is done, changes on the publisher are sent to the subscriber as they occur in real time. The subscriber applies data in the order in which commits were made on the publisher so that transactional consistency is guaranteed for the publications within any single subscription.

Logical replication is built with an architecture similar to physical streaming replication (see Section 26.2.5). It is implemented by “walsender” and “apply” processes. The walsender process starts logical decoding (described in Chapter 49) of the WAL and loads the standard logical decoding plugin (pgoutput). The plugin transforms the changes read from WAL to the logical replication protocol (see Section 53.5) and filters the data according to the publication specification. The data is then continuously transferred using the streaming replication protocol to the apply worker, which maps the data to local tables and applies the individual changes as they are received, in correct transactional order.

The apply process on the subscriber database always runs with session_replication_role set to replica, which produces the usual effects on triggers and constraints.

The logical replication apply process currently only fires row triggers, not statement triggers. The initial table synchronization, however, is implemented like a COPY command and thus fires both row and statement triggers for INSERT.

31.5.1. Initial Snapshot

The initial data in existing subscribed tables are snapshotted and copied in a parallel instance of a special kind of apply process. This process will create its own temporary replication slot and copy the existing data. Once existing data is copied, the worker enters synchronization mode, which ensures that the table is brought up to a synchronized state with the main apply process by streaming any changes that happened during the initial data copy using standard logical replication. Once the synchronization is done, the control of the replication of the table is given back to the main apply process where the replication continues as normal.

31.6. Monitoring

Because logical replication is based on a similar architecture as physical streaming replication, the monitoring on a publication node is similar to monitoring of a physical replication master (see Section 26.2.5.2).
Logical Replication

The monitoring information about subscription is visible in `pg_stat_subscription`. This view contains one row for every subscription worker. A subscription can have zero or more active subscription workers depending on its state.

Normally, there is a single apply process running for an enabled subscription. A disabled subscription or a crashed subscription will have zero rows in this view. If the initial data synchronization of any table is in progress, there will be additional workers for the tables being synchronized.

### 31.7. Security

The role used for the replication connection must have the `REPLICATION` attribute (or be a superuser). Access for the role must be configured in `pg_hba.conf` and it must have the `LOGIN` attribute.

In order to be able to copy the initial table data, the role used for the replication connection must have the `SELECT` privilege on a published table (or be a superuser).

To create a publication, the user must have the `CREATE` privilege in the database.

To add tables to a publication, the user must have ownership rights on the table. To create a publication that publishes all tables automatically, the user must be a superuser.

To create a subscription, the user must be a superuser.

The subscription apply process will run in the local database with the privileges of a superuser.

Privileges are only checked once at the start of a replication connection. They are not re-checked as each change record is read from the publisher, nor are they re-checked for each change when applied.

### 31.8. Configuration Settings

Logical replication requires several configuration options to be set.

On the publisher side, `wal_level` must be set to `logical` and `max_replication_slots` must be set to at least the number of subscriptions expected to connect, plus some reserve for table synchronization. And `max_wal_senders` should be set to at least the same as `max_replication_slots` plus the number of physical replicas that are connected at the same time.

The subscriber also requires the `max_replication_slots` to be set. In this case it should be set to at least the number of subscriptions that will be added to the subscriber. `max_logical_replication_workers` must be set to at least the number of subscriptions, again plus some reserve for the table synchronization. Additionally the `max_worker_processes` may need to be adjusted to accommodate for replication workers, at least \((\text{max\_logical\_replication\_workers} + 1)\). Note that some extensions and parallel queries also take worker slots from `max_worker_processes`.

### 31.9. Quick Setup

First set the configuration options in `postgresql.conf`:

```
wal_level = logical
```

The other required settings have default values that are sufficient for a basic setup.

`pg_hba.conf` needs to be adjusted to allow replication (the values here depend on your actual network configuration and user you want to use for connecting):

```
host     all     repuser     0.0.0.0/0     md5
```
Then on the publisher database:

```
CREATE PUBLICATION mypub FOR TABLE users, departments;
```

And on the subscriber database:

```
CREATE SUBSCRIPTION mysub CONNECTION 'dbname=foo host=bar
   user=repuser' PUBLICATION mypub;
```

The above will start the replication process, which synchronizes the initial table contents of the tables users and departments and then starts replicating incremental changes to those tables.
Chapter 32. Just-in-Time Compilation (JIT)

This chapter explains what just-in-time compilation is, and how it can be configured in PostgreSQL.

32.1. What is JIT compilation?

Just-in-Time (JIT) compilation is the process of turning some form of interpreted program evaluation into a native program, and doing so at run time. For example, instead of using general-purpose code that can evaluate arbitrary SQL expressions to evaluate a particular SQL predicate like `WHERE a.col = 3`, it is possible to generate a function that is specific to that expression and can be natively executed by the CPU, yielding a speedup.

PostgreSQL has built-in support to perform JIT compilation using LLVM\(^1\) when PostgreSQL is built with `--with-llvm`.

See `src/backend/jit/README` for further details.

32.1.1. JIT Accelerated Operations

Currently PostgreSQL's JIT implementation has support for accelerating expression evaluation and tuple deforming. Several other operations could be accelerated in the future.

Expression evaluation is used to evaluate `WHERE` clauses, target lists, aggregates and projections. It can be accelerated by generating code specific to each case.

Tuple deforming is the process of transforming an on-disk tuple (see Section 68.6.1) into its in-memory representation. It can be accelerated by creating a function specific to the table layout and the number of columns to be extracted.

32.1.2. Inlining

PostgreSQL is very extensible and allows new data types, functions, operators and other database objects to be defined; see Chapter 38. In fact the built-in objects are implemented using nearly the same mechanisms. This extensibility implies some overhead, for example due to function calls (see Section 38.3). To reduce that overhead, JIT compilation can inline the bodies of small functions into the expressions using them. That allows a significant percentage of the overhead to be optimized away.

32.1.3. Optimization

LLVM has support for optimizing generated code. Some of the optimizations are cheap enough to be performed whenever JIT is used, while others are only beneficial for longer-running queries. See https://llvm.org/docs/Passes.html#transform-passes for more details about optimizations.

32.2. When to JIT?

JIT compilation is beneficial primarily for long-running CPU-bound queries. Frequently these will be analytical queries. For short queries the added overhead of performing JIT compilation will often be higher than the time it can save.

To determine whether JIT compilation should be used, the total estimated cost of a query (see Chapter 70 and Section 19.7.2) is used. The estimated cost of the query will be compared with the setting of

\(^1\) https://llvm.org/
Just-in-Time Compilation (JIT)

jit_above_cost. If the cost is higher, JIT compilation will be performed. Two further decisions are then needed. Firstly, if the estimated cost is more than the setting of jit_inline_above_cost, short functions and operators used in the query will be inlined. Secondly, if the estimated cost is more than the setting of jit_optimize_above_cost, expensive optimizations are applied to improve the generated code. Each of these options increases the JIT compilation overhead, but can reduce query execution time considerably.

These cost-based decisions will be made at plan time, not execution time. This means that when prepared statements are in use, and a generic plan is used (see PREPARE), the values of the configuration parameters in effect at prepare time control the decisions, not the settings at execution time.

Note

If jit is set to off, or if no JIT implementation is available (for example because the server was compiled without --with-llvm), JIT will not be performed, even if it would be beneficial based on the above criteria. Setting jit to off has effects at both plan and execution time.

EXPLAIN can be used to see whether JIT is used or not. As an example, here is a query that is not using JIT:

```sql
=# EXPLAIN ANALYZE SELECT SUM(relpages) FROM pg_class;
```

<table>
<thead>
<tr>
<th>QUERY PLAN</th>
</tr>
</thead>
<tbody>
<tr>
<td>Aggregate  (cost=16.27..16.29 rows=1 width=8) (actual time=0.303..0.303 rows=1 loops=1)</td>
</tr>
<tr>
<td>-&gt;  Seq Scan on pg_class (cost=0.00..15.42 rows=342 width=4) (actual time=0.017..0.111 rows=356 loops=1)</td>
</tr>
<tr>
<td>Planning Time: 0.116 ms</td>
</tr>
<tr>
<td>Execution Time: 0.365 ms</td>
</tr>
<tr>
<td>(4 rows)</td>
</tr>
</tbody>
</table>

Given the cost of the plan, it is entirely reasonable that no JIT was used; the cost of JIT would have been bigger than the potential savings. Adjusting the cost limits will lead to JIT use:

```sql
=# SET jit_above_cost = 10;
SET
=# EXPLAIN ANALYZE SELECT SUM(relpages) FROM pg_class;
```

<table>
<thead>
<tr>
<th>QUERY PLAN</th>
</tr>
</thead>
<tbody>
<tr>
<td>Aggregate  (cost=16.27..16.29 rows=1 width=8) (actual time=6.049..6.049 rows=1 loops=1)</td>
</tr>
<tr>
<td>-&gt;  Seq Scan on pg_class (cost=0.00..15.42 rows=342 width=4) (actual time=0.019..0.052 rows=356 loops=1)</td>
</tr>
<tr>
<td>Planning Time: 0.133 ms</td>
</tr>
<tr>
<td>JIT:</td>
</tr>
<tr>
<td>Functions: 3</td>
</tr>
<tr>
<td>Options: Inlining false, Optimization false, Expressions true, Deforming true</td>
</tr>
<tr>
<td>Timing: Generation 1.259 ms, Inlining 0.000 ms, Optimization 0.797 ms, Emission 5.048 ms, Total 7.104 ms</td>
</tr>
<tr>
<td>Execution Time: 7.416 ms</td>
</tr>
</tbody>
</table>

As visible here, JIT was used, but inlining and expensive optimization were not. If jit_inline_above_cost or jit_optimize_above_cost were also lowered, that would change.
32.3. Configuration

The configuration variable jit determines whether JIT compilation is enabled or disabled. If it is enabled, the configuration variables jit_above_cost, jit_inline_above_cost, and jit_optimize_above_cost determine whether JIT compilation is performed for a query, and how much effort is spent doing so.

jit_provider determines which JIT implementation is used. It is rarely required to be changed. See Section 32.4.2.

For development and debugging purposes a few additional configuration parameters exist, as described in Section 19.17.

32.4. Extensibility

32.4.1. Inlining Support for Extensions

PostgreSQL’s JIT implementation can inline the bodies of functions of types C and internal, as well as operators based on such functions. To do so for functions in extensions, the definitions of those functions need to be made available. When using PGXS to build an extension against a server that has been compiled with LLVM JIT support, the relevant files will be built and installed automatically.

The relevant files have to be installed into $pkglibdir/bitcode/$extension/ and a summary of them into $pkglibdir/bitcode/$extension.index.bc, where $pkglibdir is the directory returned by pg_config --pkglibdir and $extension is the base name of the extension's shared library.

Note

For functions built into PostgreSQL itself, the bitcode is installed into $pkglibdir/bitcode/postgres.

32.4.2. Pluggable JIT Providers

PostgreSQL provides a JIT implementation based on LLVM. The interface to the JIT provider is pluggable and the provider can be changed without recompiling (although currently, the build process only provides inlining support data for LLVM). The active provider is chosen via the setting jit_provider.

32.4.2.1. JIT Provider Interface

A JIT provider is loaded by dynamically loading the named shared library. The normal library search path is used to locate the library. To provide the required JIT provider callbacks and to indicate that the library is actually a JIT provider, it needs to provide a C function named _PG_jit_provider_init. This function is passed a struct that needs to be filled with the callback function pointers for individual actions:

```c
struct JitProviderCallbacks
{
    JitProviderResetAfterErrorCB reset_after_error;
    JitProviderReleaseContextCB release_context;
    JitProviderCompileExprCB compile_expr;
};

extern void _PG_jit_provider_init(JitProviderCallbacks *cb);
```
Chapter 33. Regression Tests

The regression tests are a comprehensive set of tests for the SQL implementation in PostgreSQL. They test standard SQL operations as well as the extended capabilities of PostgreSQL.

33.1. Running the Tests

The regression tests can be run against an already installed and running server, or using a temporary installation within the build tree. Furthermore, there is a “parallel” and a “sequential” mode for running the tests. The sequential method runs each test script alone, while the parallel method starts up multiple server processes to run groups of tests in parallel. Parallel testing adds confidence that interprocess communication and locking are working correctly.

33.1.1. Running the Tests Against a Temporary Installation

To run the parallel regression tests after building but before installation, type:

```
make check
```

in the top-level directory. (Or you can change to `src/test/regress` and run the command there.) At the end you should see something like:

```
=======================
All 115 tests passed.
=======================
```

or otherwise a note about which tests failed. See Section 33.2 below before assuming that a “failure” represents a serious problem.

Because this test method runs a temporary server, it will not work if you did the build as the root user, since the server will not start as root. Recommended procedure is not to do the build as root, or else to perform testing after completing the installation.

If you have configured PostgreSQL to install into a location where an older PostgreSQL installation already exists, and you perform `make check` before installing the new version, you might find that the tests fail because the new programs try to use the already-installed shared libraries. (Typical symptoms are complaints about undefined symbols.) If you wish to run the tests before overwriting the old installation, you'll need to build with `configure --disable-rpath`. It is not recommended that you use this option for the final installation, however.

The parallel regression test starts quite a few processes under your user ID. Presently, the maximum concurrency is twenty parallel test scripts, which means forty processes: there's a server process and a `psql` process for each test script. So if your system enforces a per-user limit on the number of processes, make sure this limit is at least fifty or so, else you might get random-seeming failures in the parallel test. If you are not in a position to raise the limit, you can cut down the degree of parallelism by setting the `MAX_CONNECTIONS` parameter. For example:

```
make MAX_CONNECTIONS=10 check
```

runs no more than ten tests concurrently.
33.1.2. Running the Tests Against an Existing Installation

To run the tests after installation (see Chapter 16), initialize a data area and start the server as explained in Chapter 18, then type:

```
make installcheck
```

or for a parallel test:

```
make installcheck-parallel
```

The tests will expect to contact the server at the local host and the default port number, unless directed otherwise by `PGHOST` and `PGPORT` environment variables. The tests will be run in a database named `regression`; any existing database by this name will be dropped.

The tests will also transiently create some cluster-wide objects, such as roles and tablespaces. These objects will have names beginning with `regress_`. Beware of using `installcheck` mode in installations that have any actual users or tablespaces named that way.

33.1.3. Additional Test Suites

The `make check` and `make installcheck` commands run only the “core” regression tests, which test built-in functionality of the PostgreSQL server. The source distribution also contains additional test suites, most of them having to do with add-on functionality such as optional procedural languages.

To run all test suites applicable to the modules that have been selected to be built, including the core tests, type one of these commands at the top of the build tree:

```
make check-world
make installcheck-world
```

These commands run the tests using temporary servers or an already-installed server, respectively, just as previously explained for `make check` and `make installcheck`. Other considerations are the same as previously explained for each method. Note that `make check-world` builds a separate temporary installation tree for each tested module, so it requires a great deal more time and disk space than `make installcheck-world`.

Alternatively, you can run individual test suites by typing `make check` or `make installcheck` in the appropriate subdirectory of the build tree. Keep in mind that `make installcheck` assumes you’ve installed the relevant module(s), not only the core server.

The additional tests that can be invoked this way include:

- Regression tests for optional procedural languages (other than PL/pgSQL, which is tested by the core tests). These are located under `src/pl`.
- Regression tests for `contrib` modules, located under `contrib`. Not all `contrib` modules have tests.
- Regression tests for the ECPG interface library, located in `src/interfaces/ecpg/test`.
- Tests stressing behavior of concurrent sessions, located in `src/test/isolation`.
- Tests of client programs under `src/bin`. See also Section 33.4.
When using installcheck mode, these tests will destroy any existing databases named pl_regression, contrib_regression, isolation_regression, ecpg1_regression, or ecpg2_regression, as well as regression.

The TAP-based tests are run only when PostgreSQL was configured with the option --enable-tap-tests. This is recommended for development, but can be omitted if there is no suitable Perl installation.

Some test suites are not run by default, either because they are not secure to run on a multiuser system or because they require special software. You can decide which test suites to run additionally by setting the make or environment variable PG_TEST_EXTRA to a whitespace-separated list, for example:

```
make check-world PG_TEST_EXTRA='kerberos ldap ssl'
```

The following values are currently supported:

- **kerberos**
  - Runs the test suite under `src/test/kerberos`. This requires an MIT Kerberos installation and opens TCP/IP listen sockets.

- **ldap**
  - Runs the test suite under `src/test/ldap`. This requires an OpenLDAP installation and opens TCP/IP listen sockets.

- **ssl**
  - Runs the test suite under `src/test/ssl`. This opens TCP/IP listen sockets.

Tests for features that are not supported by the current build configuration are not run even if they are mentioned in PG_TEST_EXTRA.

### 33.1.4. Locale and Encoding

By default, tests using a temporary installation use the locale defined in the current environment and the corresponding database encoding as determined by `initdb`. It can be useful to test different locales by setting the appropriate environment variables, for example:

```
make check LANG=C
make check LC_COLLATE=en_US.utf8 LC_CTYPE=fr_CA.utf8
```

For implementation reasons, setting LC_ALL does not work for this purpose; all the other locale-related environment variables do work.

When testing against an existing installation, the locale is determined by the existing database cluster and cannot be set separately for the test run.

You can also choose the database encoding explicitly by setting the variable ENCODING, for example:

```
make check LANG=C ENCODING=EUC_JP
```

Setting the database encoding this way typically only makes sense if the locale is C; otherwise the encoding is chosen automatically from the locale, and specifying an encoding that does not match the locale will result in an error.

The database encoding can be set for tests against either a temporary or an existing installation, though in the latter case it must be compatible with the installation's locale.
33.1.5. Extra Tests

The core regression test suite contains a few test files that are not run by default, because they might be platform-dependent or take a very long time to run. You can run these or other extra test files by setting the variable EXTRA_TESTS. For example, to run the numeric_big test:

```
make check EXTRA_TESTS=numeric_big
```

To run the collation tests:

```
make check EXTRA_TESTS='collate.icu.utf8 collate.linux.utf8'
LANG=en_US.utf8
```

The collate.linux.utf8 test works only on Linux/glibc platforms. The collate.icu.utf8 test only works when support for ICU was built. Both tests will only succeed when run in a database that uses UTF-8 encoding.

33.1.6. Testing Hot Standby

The source distribution also contains regression tests for the static behavior of Hot Standby. These tests require a running primary server and a running standby server that is accepting new WAL changes from the primary (using either file-based log shipping or streaming replication). Those servers are not automatically created for you, nor is replication setup documented here. Please check the various sections of the documentation devoted to the required commands and related issues.

To run the Hot Standby tests, first create a database called regression on the primary:

```
psql -h primary -c "CREATE DATABASE regression"
```

Next, run the preparatory script src/test/regress/sql/hs_primary_setup.sql on the primary in the regression database, for example:

```
psql -h primary -f src/test/regress/sql/hs_primary_setup.sql regression
```

Allow these changes to propagate to the standby.

Now arrange for the default database connection to be to the standby server under test (for example, by setting the PGHOST and PGPORT environment variables). Finally, run make standbycheck in the regression directory:

```
cd src/test/regress
make standbycheck
```

Some extreme behaviors can also be generated on the primary using the script src/test/regress/sql/hs_primary_extremes.sql to allow the behavior of the standby to be tested.

33.2. Test Evaluation

Some properly installed and fully functional PostgreSQL installations can “fail” some of these regression tests due to platform-specific artifacts such as varying floating-point representation and message wording. The tests are currently evaluated using a simple diff comparison against the outputs generated on a reference system, so the results are sensitive to small system differences. When a test is
Regression Tests

reported as “failed”, always examine the differences between expected and actual results; you might
find that the differences are not significant. Nonetheless, we still strive to maintain accurate reference
files across all supported platforms, so it can be expected that all tests pass.

The actual outputs of the regression tests are in files in the `src/test/regress/results` direc-
tory. The test script uses `diff` to compare each output file against the reference outputs stored in
the `src/test/regress/expected` directory. Any differences are saved for your inspection in
`src/test/regress/regression.diffs`. (When running a test suite other than the core tests,
these files of course appear in the relevant subdirectory, not `src/test/regress`.)

If you don't like the `diff` options that are used by default, set the environment variable `PG_RE-
GRESS_DIFF_OPTS`, for instance `PG_REGRESS_DIFF_OPTS='-u'`. (Or you can run `diff`
yourself, if you prefer.)

If for some reason a particular platform generates a “failure” for a given test, but inspection of the
output convinces you that the result is valid, you can add a new comparison file to silence the failure
report in future test runs. See Section 33.3 for details.

### 33.2.1. Error Message Differences

Some of the regression tests involve intentional invalid input values. Error messages can come from
either the PostgreSQL code or from the host platform system routines. In the latter case, the messages
can vary between platforms, but should reflect similar information. These differences in messages will
result in a “failed” regression test that can be validated by inspection.

### 33.2.2. Locale Differences

If you run the tests against a server that was initialized with a collation-order locale other than C, then
there might be differences due to sort order and subsequent failures. The regression test suite is set
up to handle this problem by providing alternate result files that together are known to handle a large
number of locales.

To run the tests in a different locale when using the temporary-installation method, pass the appropriate
locale-related environment variables on the `make` command line, for example:

```
make check LANG=de_DE.utf8
```

(The regression test driver unsets `LC_ALL`, so it does not work to choose the locale using that variable.)

To use no locale, either unset all locale-related environment variables (or set them to `C`) or use the
following special invocation:

```
make check NO_LOCALE=1
```

When running the tests against an existing installation, the locale setup is determined by the existing
installation. To change it, initialize the database cluster with a different locale by passing the appro-
priate options to `initdb`.

In general, it is advisable to try to run the regression tests in the locale setup that is wanted for pro-
duction use, as this will exercise the locale- and encoding-related code portions that will actually be
used in production. Depending on the operating system environment, you might get failures, but then
you will at least know what locale-specific behaviors to expect when running real applications.

### 33.2.3. Date and Time Differences

Most of the date and time results are dependent on the time zone environment. The reference files are
generated for time zone `PST8PDT` (Berkeley, California), and there will be apparent failures if the
tests are not run with that time zone setting. The regression test driver sets environment variable PGTZ to PST8PDT, which normally ensures proper results.

33.2.4. Floating-Point Differences

Some of the tests involve computing 64-bit floating-point numbers (double precision) from table columns. Differences in results involving mathematical functions of double precision columns have been observed. The float8 and geometry tests are particularly prone to small differences across platforms, or even with different compiler optimization settings. Human eyeball comparison is needed to determine the real significance of these differences which are usually 10 places to the right of the decimal point.

Some systems display minus zero as -0, while others just show 0.

Some systems signal errors from pow() and exp() differently from the mechanism expected by the current PostgreSQL code.

33.2.5. Row Ordering Differences

You might see differences in which the same rows are output in a different order than what appears in the expected file. In most cases this is not, strictly speaking, a bug. Most of the regression test scripts are not so pedantic as to use an ORDER BY for every single SELECT, and so their result row orderings are not well-defined according to the SQL specification. In practice, since we are looking at the same queries being executed on the same data by the same software, we usually get the same result ordering on all platforms, so the lack of ORDER BY is not a problem. Some queries do exhibit cross-platform ordering differences, however. When testing against an already-installed server, ordering differences can also be caused by non-C locale settings or non-default parameter settings, such as custom values of work_mem or the planner cost parameters.

Therefore, if you see an ordering difference, it’s not something to worry about, unless the query does have an ORDER BY that your result is violating. However, please report it anyway, so that we can add an ORDER BY to that particular query to eliminate the bogus “failure” in future releases.

You might wonder why we don’t order all the regression test queries explicitly to get rid of this issue once and for all. The reason is that that would make the regression tests less useful, not more, since they’d tend to exercise query plan types that produce ordered results to the exclusion of those that don’t.

33.2.6. Insufficient Stack Depth

If the errors test results in a server crash at the select infinite_recurse() command, it means that the platform’s limit on process stack size is smaller than the max_stack_depth parameter indicates. This can be fixed by running the server under a higher stack size limit (4MB is recommended with the default value of max_stack_depth). If you are unable to do that, an alternative is to reduce the value of max_stack_depth.

On platforms supporting getrlimit(), the server should automatically choose a safe value of max_stack_depth; so unless you’ve manually overridden this setting, a failure of this kind is a reportable bug.

33.2.7. The “random” Test

The random test script is intended to produce random results. In very rare cases, this causes that regression test to fail. Typing:

diff results/random.out expected/random.out

should produce only one or a few lines of differences. You need not worry unless the random test fails repeatedly.
33.2.8. Configuration Parameters

When running the tests against an existing installation, some non-default parameter settings could cause the tests to fail. For example, changing parameters such as `enable_seqscan` or `enable_indexscan` could cause plan changes that would affect the results of tests that use `EXPLAIN`.

33.3. Variant Comparison Files

Since some of the tests inherently produce environment-dependent results, we have provided ways to specify alternate “expected” result files. Each regression test can have several comparison files showing possible results on different platforms. There are two independent mechanisms for determining which comparison file is used for each test.

The first mechanism allows comparison files to be selected for specific platforms. There is a mapping file, `src/test/regress/resultmap`, that defines which comparison file to use for each platform. To eliminate bogus test “failures” for a particular platform, you first choose or make a variant result file, and then add a line to the `resultmap` file.

Each line in the mapping file is of the form

```
testname:output:platformpattern=comparisonfilename
```

The test name is just the name of the particular regression test module. The output value indicates which output file to check. For the standard regression tests, this is always `out`. The value corresponds to the file extension of the output file. The platform pattern is a pattern in the style of the Unix tool `expr` (that is, a regular expression with an implicit `^` anchor at the start). It is matched against the platform name as printed by `config.guess`. The comparison file name is the base name of the substitute result comparison file.

For example: some systems interpret very small floating-point values as zero, rather than reporting an underflow error. This causes a few differences in the `float8` regression test. Therefore, we provide a variant comparison file, `float8-small-is-zero.out`, which includes the results to be expected on these systems. To silence the bogus “failure” message on OpenBSD platforms, `resultmap` includes:

```
float8:out:i.86-.*-openbsd=float8-small-is-zero.out
```

which will trigger on any machine where the output of `config.guess` matches `i.86-.*-openbsd`. Other lines in `resultmap` select the variant comparison file for other platforms where it’s appropriate.

The second selection mechanism for variant comparison files is much more automatic: it simply uses the “best match” among several supplied comparison files. The regression test driver script considers both the standard comparison file for a test, `testname.out`, and variant files named `testname_digit.out` (where the `digit` is any single digit `0-9`). If any such file is an exact match, the test is considered to pass; otherwise, the one that generates the shortest diff is used to create the failure report. (If `resultmap` includes an entry for the particular test, then the base `testname` is the substitute name given in `resultmap`.)

For example, for the `char` test, the comparison file `char.out` contains results that are expected in the C and POSIX locales, while the file `char_1.out` contains results sorted as they appear in many other locales.

The best-match mechanism was devised to cope with locale-dependent results, but it can be used in any situation where the test results cannot be predicted easily from the platform name alone. A limitation of this mechanism is that the test driver cannot tell which variant is actually “correct” for the current
Regression Tests

environment; it will just pick the variant that seems to work best. Therefore it is safest to use this mechanism only for variant results that you are willing to consider equally valid in all contexts.

33.4. TAP Tests

Various tests, particularly the client program tests under src/bin, use the Perl TAP tools and are run using the Perl testing program prove. You can pass command-line options to prove by setting the make variable PROVE_FLAGS, for example:

```
make -C src/bin check PROVE_FLAGS='--timer'
```

See the manual page of prove for more information.

The make variable PROVE_TESTS can be used to define a whitespace-separated list of paths relative to the Makefile invoking prove to run the specified subset of tests instead of the default t/*.pl. For example:

```
make check PROVE_TESTS='t/001_test1.pl t/003_test3.pl'
```

The TAP tests require the Perl module IPC::Run. This module is available from CPAN or an operating system package.

33.5. Test Coverage Examination

The PostgreSQL source code can be compiled with coverage testing instrumentation, so that it becomes possible to examine which parts of the code are covered by the regression tests or any other test suite that is run with the code. This is currently supported when compiling with GCC and requires the gcov and lcov programs.

A typical workflow would look like this:

```
./configure --enable-coverage ... OTHER OPTIONS ...
make
make check # or other test suite
make coverage-html
```

Then point your HTML browser to coverage/index.html. The make commands also work in subdirectories.

If you don't have lcov or prefer text output over an HTML report, you can also run

```
make coverage
```

instead of make coverage-html, which will produce .gcov output files for each source file relevant to the test. (make coverage and make coverage-html will overwrite each other's files, so mixing them might be confusing.)

To reset the execution counts between test runs, run:

```
make coverage-clean
```
Part IV. Client Interfaces

This part describes the client programming interfaces distributed with PostgreSQL. Each of these chapters can be read independently. Note that there are many other programming interfaces for client programs that are distributed separately and contain their own documentation (Appendix H lists some of the more popular ones). Readers of this part should be familiar with using SQL commands to manipulate and query the database (see Part II) and of course with the programming language that the interface uses.
# Table of Contents

34. libpq - C Library ................................................................. 770

34.1. Database Connection Control Functions .................................. 770
  34.1.1. Connection Strings .................................................. 777
  34.1.2. Parameter Key Words ................................................ 778
34.2. Connection Status Functions ............................................. 783
34.3. Command Execution Functions ........................................... 789
  34.3.1. Main Functions ....................................................... 789
  34.3.2. Retrieving Query Result Information ................................ 797
  34.3.3. Retrieving Other Result Information ................................ 801
  34.3.4. Escaping Strings for Inclusion in SQL Commands ................. 802
34.4. Asynchronous Command Processing ....................................... 805
34.5. Retrieving Query Results Row-By-Row .................................. 809
34.6. Canceling Queries in Progress .......................................... 810
34.7. The Fast-Path Interface .................................................. 811
34.8. Asynchronous Notification .............................................. 812
34.9. Functions Associated with the COPY Command ......................... 813
  34.9.1. Functions for Sending COPY Data .................................. 814
  34.9.2. Functions for Receiving COPY Data ................................ 814
  34.9.3. Obsolete Functions for COPY ....................................... 815
34.10. Control Functions ........................................................ 817
34.11. Miscellaneous Functions ................................................ 819
34.12. Notice Processing ........................................................ 822
34.13. Event System ............................................................. 823
  34.13.1. Event Types .......................................................... 823
  34.13.2. Event Callback Procedure ......................................... 825
  34.13.3. Event Support Functions ......................................... 826
  34.13.4. Event Example ...................................................... 827
34.14. Environment Variables .................................................. 829
34.15. The Password File ....................................................... 831
34.16. The Connection Service File .......................................... 831
34.17. LDAP Lookup of Connection Parameters ................................. 832
34.18. SSL Support .............................................................. 833
  34.18.1. Client Verification of Server Certificates ......................... 833
  34.18.2. Client Certificates ................................................ 834
  34.18.3. Protection Provided in Different Modes ........................... 834
  34.18.4. SSL Client File Usage ........................................... 835
  34.18.5. SSL Library Initialization ....................................... 836
34.19. Behavior in Threaded Programs ......................................... 836
34.20. Building libpq Programs ................................................. 837
34.21. Example Programs ........................................................ 838

35. Large Objects ............................................................... 850

35.1. Introduction ..................................................................... 850
35.2. Implementation Features ................................................ 850
35.3. Client Interfaces .......................................................... 850
  35.3.1. Creating a Large Object ............................................. 851
  35.3.2. Importing a Large Object ........................................... 851
  35.3.3. Exporting a Large Object ........................................... 852
  35.3.4. Opening an Existing Large Object .................................. 852
  35.3.5. Writing Data to a Large Object .................................... 852
  35.3.6. Reading Data from a Large Object ................................ 853
  35.3.7. Seeking in a Large Object ......................................... 853
  35.3.8. Obtaining the Seek Position of a Large Object .................. 853
  35.3.9. Truncating a Large Object ......................................... 854
  35.3.10. Closing a Large Object Descriptor ................................. 854
  35.3.11. Removing a Large Object ......................................... 854

766
36. ECPG - Embedded SQL in C ...................................................... 862
36.1. The Concept ................................................................. 862
36.2. Managing Database Connections .................................... 862
  36.2.1. Connecting to the Database Server ............................ 862
  36.2.2. Choosing a Connection ............................................ 864
  36.2.3. Closing a Connection ............................................. 865
36.3. Running SQL Commands ................................................ 865
  36.3.1. Executing SQL Statements ...................................... 865
  36.3.2. Using Cursors .................................................... 866
  36.3.3. Managing Transactions ........................................ 867
  36.3.4. Prepared Statements ........................................... 867
36.4. Using Host Variables .................................................. 868
  36.4.1. Overview ............................................................ 868
  36.4.2. Declare Sections .................................................. 868
  36.4.3. Retrieving Query Results ..................................... 869
  36.4.4. Type Mapping ..................................................... 870
  36.4.5. Handling Nonprimitive SQL Data Types ..................... 877
  36.4.6. Indicators .......................................................... 881
36.5. Dynamic SQL ................................................................ 882
  36.5.1. Executing Statements without a Result Set ............... 882
  36.5.2. Executing a Statement with Input Parameters ............. 882
  36.5.3. Executing a Statement with a Result Set .................. 883
36.6. pgtypes Library ............................................................ 884
  36.6.1. Character Strings .................................................. 884
  36.6.2. The numeric Type .................................................. 884
  36.6.3. The date Type ...................................................... 887
  36.6.4. The timestamp Type .............................................. 891
  36.6.5. The interval Type ................................................ 895
  36.6.6. The decimal Type ............................................... 896
  36.6.7. errno Values of pgtypeslib .................................... 896
  36.6.8. Special Constants of pgtypeslib ............................... 897
36.7. Using Descriptor Areas .................................................. 898
  36.7.1. Named SQL Descriptor Areas ................................. 898
  36.7.2. SQLDA Descriptor Areas ....................................... 900
36.8. Error Handling ............................................................ 911
  36.8.1. Setting Callbacks .................................................. 911
  36.8.2. sqlca ............................................................... 913
  36.8.3. SQLSTATE vs. SQLCODE ...................................... 914
36.9. Preprocessor Directives ............................................... 918
  36.9.1. Including Files ................................................... 918
  36.9.2. The define and undef Directives ............................. 919
  36.9.3. ifdef, ifndef, else, elif, and endif Directives .......... 919
36.10. Processing Embedded SQL Programs ................................. 920
36.11. Library Functions ........................................................ 921
36.12. Large Objects ............................................................. 922
36.13. C++ Applications ......................................................... 923
  36.13.1. Scope for Host Variables .............................. 924
  36.13.2. C++ Application Development with External C Module 925
36.14. Embedded SQL Commands .............................................. 927
36.15. Informix Compatibility Mode ......................................... 951
  36.15.1. Additional Types ............................................... 951
  36.15.2. Additional/Missing Embedded SQL Statements ............ 952
  36.15.3. Informix-compatible SQLDA Descriptor Areas .......... 952
  36.15.4. Additional Functions ......................................... 955
  36.15.5. Additional Constants ......................................... 965
36.16. Internals ................................................................. 966
37. The Information Schema ................................................................................................................. 969
    37.1. The Schema .............................................................................................................................. 969
    37.2. Data Types ................................................................................................................................ 969
    37.3. information_schema_catalog_name ................................................................. 970
    37.4. administrable_role_authorizations .......................................................... 970
    37.5. applicable_roles ......................................................... 970
    37.6. attributes ................................................................................................................................. 971
    37.7. character_sets ......................................................................................................................... 974
    37.8. check_constraint_routine_usage ......................................................... 975
    37.9. check_constraints .................................................................................................................... 975
    37.10. collations ............................................................................................................................... 976
    37.11. collation_character_set_applicability .............................................................. 976
    37.12. column_domain_usage ......................................................................................................... 977
    37.13. column_options ...................................................................................................................... 977
    37.14. column_privileges .................................................................................................................. 977
    37.15. column_udt_usage ................................................................................................................... 978
    37.16. columns .................................................................................................................................. 979
    37.17. constraint_column_usage ....................................................................................................... 983
    37.18. constraint_table_usage .......................................................................................................... 983
    37.19. data_type_privileges .............................................................................................................. 984
    37.20. domain_constraints ............................................................................................................... 985
    37.21. domain_udt_usage .................................................................................................................... 985
    37.22. domains .................................................................................................................................. 986
    37.23. element_types ......................................................................................................................... 988
    37.24. enabled_roles .......................................................................................................................... 991
    37.25. foreign_data_wrapper_options ............................................................................................... 991
    37.26. foreign_data_wrappers ........................................................................................................... 991
    37.27. foreign_server_options ............................................................................................................ 992
    37.28. foreign_servers ....................................................................................................................... 992
    37.29. foreign_table_options ............................................................................................................. 993
    37.30. foreign_tables .......................................................................................................................... 993
    37.31. key_column_usage .................................................................................................................. 993
    37.32. parameters .............................................................................................................................. 994
    37.33. referential_constraints .......................................................................................................... 996
    37.34. role_column_grants ................................................................................................................. 997
    37.35. role_routine_grants ................................................................................................................ 998
    37.36. role_table_grants .................................................................................................................... 998
    37.37. role_udt_grants ....................................................................................................................... 999
    37.38. role_usage_grants ................................................................................................................... 1000
    37.39. routine_privileges .................................................................................................................... 1000
    37.40. routines .................................................................................................................................. 1001
    37.41. schemata ................................................................................................................................. 1006
    37.42. sequences ............................................................................................................................... 1006
    37.43. sql_features ............................................................................................................................. 1007
    37.44. sql_impementation_info .......................................................................................................... 1008
    37.45. sql_languages .......................................................................................................................... 1008
    37.46. sql_packages ........................................................................................................................... 1009
    37.47. sql_parts ................................................................................................................................. 1009
    37.48. sql_sizing ................................................................................................................................. 1010
    37.49. sql_sizing_profiles .................................................................................................................. 1010
    37.50. table_constraints .................................................................................................................... 1011
    37.51. table_privileges ....................................................................................................................... 1011
    37.52. tables ..................................................................................................................................... 1012
    37.53. transforms ............................................................................................................................... 1013
    37.54. triggered_update_columns .................................................................................................... 1014
    37.55. triggers .................................................................................................................................... 1014
    37.56. udt_privileges .......................................................................................................................... 1016
    37.57. usage_privileges ...................................................................................................................... 1016
<table>
<thead>
<tr>
<th>Section</th>
<th>Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>37.58.</td>
<td>user_defined_types</td>
<td>1017</td>
</tr>
<tr>
<td>37.59.</td>
<td>user_mapping_options</td>
<td>1018</td>
</tr>
<tr>
<td>37.60.</td>
<td>user_mappings</td>
<td>1019</td>
</tr>
<tr>
<td>37.61.</td>
<td>view_column_usage</td>
<td>1019</td>
</tr>
<tr>
<td>37.62.</td>
<td>view_routine_usage</td>
<td>1020</td>
</tr>
<tr>
<td>37.63.</td>
<td>view_table_usage</td>
<td>1020</td>
</tr>
<tr>
<td>37.64.</td>
<td>views</td>
<td>1021</td>
</tr>
</tbody>
</table>
Chapter 34. libpq - C Library

libpq is the C application programmer's interface to PostgreSQL. libpq is a set of library functions that allow client programs to pass queries to the PostgreSQL backend server and to receive the results of these queries.

libpq is also the underlying engine for several other PostgreSQL application interfaces, including those written for C++, Perl, Python, Tcl and ECPG. So some aspects of libpq's behavior will be important to you if you use one of those packages. In particular, Section 34.14, Section 34.15 and Section 34.18 describe behavior that is visible to the user of any application that uses libpq.

Some short programs are included at the end of this chapter (Section 34.21) to show how to write programs that use libpq. There are also several complete examples of libpq applications in the directory src/test/examples in the source code distribution.

Client programs that use libpq must include the header file libpq-fe.h and must link with the libpq library.

34.1. Database Connection Control Functions

The following functions deal with making a connection to a PostgreSQL backend server. An application program can have several backend connections open at one time. (One reason to do that is to access more than one database.) Each connection is represented by a PGconn object, which is obtained from the function PQconnectdb, PQconnectdbParams, or PQsetdbLogin. Note that these functions will always return a non-null object pointer, unless perhaps there is too little memory even to allocate the PGconn object. The PQstatus function should be called to check the return value for a successful connection before queries are sent via the connection object.

**Warning**

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from search_path. One can set parameter key word options to value -c search_path=. Alternately, one can issue PQexec(conn, "SELECT pg_catalog.set_config('search_path', '', false)") after connecting. This consideration is not specific to libpq; it applies to every interface for executing arbitrary SQL commands.

**Warning**

On Unix, forking a process with open libpq connections can lead to unpredictable results because the parent and child processes share the same sockets and operating system resources. For this reason, such usage is not recommended, though doing an exec from the child process to load a new executable is safe.

**Note**

On Windows, there is a way to improve performance if a single database connection is repeatedly started and shutdown. Internally, libpq calls WSAStartup() and WSACleanup() for connection startup and shutdown, respectively. WSAStartup() increments an internal Windows library reference count which is decremented by WSA-
Cleanup(). When the reference count is just one, calling WSACleanup() frees all resources and all DLLs are unloaded. This is an expensive operation. To avoid this, an application can manually call WSAStrartup() so resources will not be freed when the last database connection is closed.

**PQconnectdbParams**

Makes a new connection to the database server.

```c
PGconn *PQconnectdbParams(const char * const *keywords,
    const char * const *values,
    int expand_dbname);
```

This function opens a new database connection using the parameters taken from two NULL-terminated arrays. The first, **keywords**, is defined as an array of strings, each one being a key word. The second, **values**, gives the value for each key word. Unlike PQsetdbLogin below, the parameter set can be extended without changing the function signature, so use of this function (or its nonblocking analogs PQconnectStartParams and PQconnectPoll) is preferred for new application programming.

The currently recognized parameter key words are listed in Section 34.1.2.

When expand_dbname is non-zero, the **dbname** key word value is allowed to be recognized as a connection string. Only the first occurrence of **dbname** is expanded this way, any subsequent **dbname** value is processed as plain database name. More details on the possible connection string formats appear in Section 34.1.1.

The passed arrays can be empty to use all default parameters, or can contain one or more parameter settings. They should be matched in length. Processing will stop at the first NULL element in the **keywords** array.

If any parameter is NULL or an empty string, the corresponding environment variable (see Section 34.14) is checked. If the environment variable is not set either, then the indicated built-in defaults are used.

In general key words are processed from the beginning of these arrays in index order. The effect of this is that when key words are repeated, the last processed value is retained. Therefore, through careful placement of the **dbname** key word, it is possible to determine what may be overridden by a **conninfo** string, and what may not.

**PQconnectdb**

Makes a new connection to the database server.

```c
PGconn *PQconnectdb(const char *conninfo);
```

This function opens a new database connection using the parameters taken from the string **conninfo**.

The passed string can be empty to use all default parameters, or it can contain one or more parameter settings separated by whitespace, or it can contain a URI. See Section 34.1.1 for details.

**PQsetdbLogin**

Makes a new connection to the database server.
libpq - C Library

PGconn *PQsetdbLogin(const char *pghost,
        const char *pgport,
        const char *pgoptions,
        const char *pgtty,
        const char *dbName,
        const char *login,
        const char *pwd);

This is the predecessor of PQconnectdb with a fixed set of parameters. It has the same functionality except that the missing parameters will always take on default values. Write NULL or an empty string for any one of the fixed parameters that is to be defaulted.

If the dbName contains an = sign or has a valid connection URI prefix, it is taken as a conninfo string in exactly the same way as if it had been passed to PQconnectdb, and the remaining parameters are then applied as specified for PQconnectdbParams.

PQsetdb

Makes a new connection to the database server.

PGconn *PQsetdb(char *pghost,
        char *pgport,
        char *pgoptions,
        char *pgtty,
        char *dbName);

This is a macro that calls PQsetdbLogin with null pointers for the login and pwd parameters. It is provided for backward compatibility with very old programs.

PQconnectStartParams
PQconnectStart
PQconnectPoll

Make a connection to the database server in a nonblocking manner.

PGconn *PQconnectStartParams(const char * const *keywords,
        const char * const *values,
        int expand_dbname);

PGconn *PQconnectStart(const char *conninfo);

PostgresPollingStatusType PQconnectPoll(PGconn *conn);

These three functions are used to open a connection to a database server such that your application's thread of execution is not blocked on remote I/O whilst doing so. The point of this approach is that the waits for I/O to complete can occur in the application's main loop, rather than down inside PQconnectdbParams or PQconnectdb, and so the application can manage this operation in parallel with other activities.

With PQconnectStartParams, the database connection is made using the parameters taken from the keywords and values arrays, and controlled by expand_dbname, as described above for PQconnectdbParams.

With PQconnectStart, the database connection is made using the parameters taken from the string conninfo as described above for PQconnectdb.

Neither PQconnectStartParams nor PQconnectStart nor PQconnectPoll will block, so long as a number of restrictions are met:
• The hostaddr parameter must be used appropriately to prevent DNS queries from being made. See the documentation of this parameter in Section 34.1.2 for details.

• If you call PQtrace, ensure that the stream object into which you trace will not block.

• You must ensure that the socket is in the appropriate state before calling PQconnectPoll, as described below.

To begin a nonblocking connection request, call PQconnectStart or PQconnectStartParams. If the result is null, then libpq has been unable to allocate a new PGconn structure. Otherwise, a valid PGconn pointer is returned (though not yet representing a valid connection to the database). Next call PQstatus(conn). If the result is CONNECTION_BAD, the connection attempt has already failed, typically because of invalid connection parameters.

If PQconnectStart or PQconnectStartParams succeeds, the next stage is to poll libpq so that it can proceed with the connection sequence. Use PQsocket(conn) to obtain the descriptor of the socket underlying the database connection. (Caution: do not assume that the socket remains the same across PQconnectPoll calls.) Loop thus: If PQconnectPoll(conn) last returned PGRES_POLLING_READING, wait until the socket is ready to read (as indicated by select(), poll(), or similar system function). Then call PQconnectPoll(conn) again. Conversely, if PQconnectPoll(conn) last returned PGRES_POLLING_WRITING, wait until the socket is ready to write, then call PQconnectPoll(conn) again. On the first iteration, i.e. if you have yet to call PQconnectPoll, behave as if it last returned PGRES_POLLING_WRITING. Continue this loop until PQconnectPoll(conn) returns PGRES_POLLING_FAILED, indicating the connection procedure has failed, or PGRES_POLLING_OK, indicating the connection has been successfully made.

At any time during connection, the status of the connection can be checked by calling PQstatus. If this call returns CONNECTION_BAD, then the connection procedure has failed; if the call returns CONNECTION_OK, then the connection is ready. Both of these states are equally detectable from the return value of PQconnectPoll, described above. Other states might also occur during (and only during) an asynchronous connection procedure. These indicate the current stage of the connection procedure and might be useful to provide feedback to the user for example. These statuses are:

CONNECTION_STARTED

Waiting for connection to be made.

CONNECTION_MADE

Connection OK; waiting to send.

CONNECTION_WAITING_RESPONSE

Waiting for a response from the server.

CONNECTION_AUTH_OK

Received authentication; waiting for backend start-up to finish.

CONNECTION_SSL_STARTUP

Negotiating SSL encryption.

CONNECTION_SETENV

Negotiating environment-driven parameter settings.

CONNECTION_CHECK_WRITABLE

Checking if connection is able to handle write transactions.
CONNECTION_CONSUME

Consuming any remaining response messages on connection.

Note that, although these constants will remain (in order to maintain compatibility), an application should never rely upon these occurring in a particular order, or at all, or on the status always being one of these documented values. An application might do something like this:

```c
switch(PQstatus(conn))
{
    case CONNECTION_STARTED:
        feedback = "Connecting...";
        break;

    case CONNECTION_MADE:
        feedback = "Connected to server...";
        break;

    default:
        feedback = "Connecting...";
}
```

The connect_timeout connection parameter is ignored when using PQconnectPoll; it is the application's responsibility to decide whether an excessive amount of time has elapsed. Otherwise, PQconnectStart followed by a PQconnectPoll loop is equivalent to PQconnectdb.

Note that when PQconnectStart or PQconnectStartParams returns a non-null pointer, you must call PQfinish when you are finished with it, in order to dispose of the structure and any associated memory blocks. This must be done even if the connection attempt fails or is abandoned.

PQconndefaults

Returns the default connection options.

```c
PQconninfoOption *PQconndefaults(void);
```

typedef struct
{
    char *keyword;   /* The keyword of the option */
    char *envvar;    /* Fallback environment variable name */
    char *compiled;  /* Fallback compiled in default value */
    char *val;       /* Option's current value, or NULL */
    char *label;     /* Label for field in connect dialog */
    char *dispchar;  /* Indicates how to display this field in a connect dialog. Values are:
                      
                      ""        Display entered value as is
                      "*"       Password field - hide value
                      "D"       Debug option - don't show by default */
    int dispsize;    /* Field size in characters for dialog */

} PQconninfoOption;

Returns a connection options array. This can be used to determine all possible PQconnectdb options and their current default values. The return value points to an array of PQconninfoOp-
tion structures, which ends with an entry having a null keyword pointer. The null pointer is returned if memory could not be allocated. Note that the current default values (val fields) will depend on environment variables and other context. A missing or invalid service file will be silently ignored. Callers must treat the connection options data as read-only.

After processing the options array, free it by passing it to PQconninfoFree. If this is not done, a small amount of memory is leaked for each call to PQconndefaults.

**PQconninfo**

Returns the connection options used by a live connection.

```c
PQconninfoOption *PQconninfo(PGconn *conn);
```

Returns a connection options array. This can be used to determine all possible PQconnectdb options and the values that were used to connect to the server. The return value points to an array of PQconninfoOption structures, which ends with an entry having a null keyword pointer. All notes above for PQconndefaults also apply to the result of PQconninfo.

**PQconninfoParse**

Returns parsed connection options from the provided connection string.

```c
PQconninfoOption *PQconninfoParse(const char *conninfo, char **errmsg);
```

Parses a connection string and returns the resulting options as an array; or returns NULL if there is a problem with the connection string. This function can be used to extract the PQconnectdb options in the provided connection string. The return value points to an array of PQconninfoOption structures, which ends with an entry having a null keyword pointer.

All legal options will be present in the result array, but the PQconninfoOption for any option not present in the connection string will have val set to NULL; default values are not inserted.

If errmsg is not NULL, then *errmsg is set to NULL on success, else to a malloc’d error string explaining the problem. (It is also possible for *errmsg to be set to NULL and the function to return NULL; this indicates an out-of-memory condition.)

After processing the options array, free it by passing it to PQconninfoFree. If this is not done, some memory is leaked for each call to PQconninfoParse. Conversely, if an error occurs and errmsg is not NULL, be sure to free the error string using PQfreemem.

**PQfinish**

Closes the connection to the server. Also frees memory used by the PGconn object.

```c
void PQfinish(PGconn *conn);
```

Note that even if the server connection attempt fails (as indicated by PQstatus), the application should call PQfinish to free the memory used by the PGconn object. The PGconn pointer must not be used again after PQfinish has been called.

**PQreset**

Resets the communication channel to the server.
void PQreset(PGconn *conn);

This function will close the connection to the server and attempt to reestablish a new connection to the same server, using all the same parameters previously used. This might be useful for error recovery if a working connection is lost.

PQresetStart
PQresetPoll

Reset the communication channel to the server, in a nonblocking manner.

int PQresetStart(PGconn *conn);

PostgresPollingStatusType PQresetPoll(PGconn *conn);

These functions will close the connection to the server and attempt to reestablish a new connection to the same server, using all the same parameters previously used. This can be useful for error recovery if a working connection is lost. They differ from PQreset (above) in that they act in a nonblocking manner. These functions suffer from the same restrictions as PQconnectStartParams, PQconnectStart and PQconnectPoll.

To initiate a connection reset, call PQresetStart. If it returns 0, the reset has failed. If it returns 1, poll the reset using PQresetPoll in exactly the same way as you would create the connection using PQconnectPoll.

PQpingParams

PQpingParams reports the status of the server. It accepts connection parameters identical to those of PQconnectdbParams, described above. It is not necessary to supply correct user name, password, or database name values to obtain the server status; however, if incorrect values are provided, the server will log a failed connection attempt.

PGPing PQpingParams(const char * const *keywords, const char * const *values, int expand_dbname);

The function returns one of the following values:

PQPING_OK

The server is running and appears to be accepting connections.

PQPING_REJECT

The server is running but is in a state that disallows connections (startup, shutdown, or crash recovery).

PQPING_NO_RESPONSE

The server could not be contacted. This might indicate that the server is not running, or that there is something wrong with the given connection parameters (for example, wrong port number), or that there is a network connectivity problem (for example, a firewall blocking the connection request).

PQPING_NO_ATTEMPT

No attempt was made to contact the server, because the supplied parameters were obviously incorrect or there was some client-side problem (for example, out of memory).
PQping

PQping reports the status of the server. It accepts connection parameters identical to those of PQconnectdb, described above. It is not necessary to supply correct user name, password, or database name values to obtain the server status; however, if incorrect values are provided, the server will log a failed connection attempt.

PGPing PQping(const char *conninfo);

The return values are the same as for PQpingParams.

34.1.1. Connection Strings

Several libpq functions parse a user-specified string to obtain connection parameters. There are two accepted formats for these strings: plain keyword = value strings and URIs. URIs generally follow RFC 39861, except that multi-host connection strings are allowed as further described below.

34.1.1.1. Keyword/Value Connection Strings

In the first format, each parameter setting is in the form keyword = value. Spaces around the equal sign are optional. To write an empty value, or a value containing spaces, surround it with single quotes, e.g., keyword = 'a value'. Single quotes and backslashes within the value must be escaped with a backslash, i.e., \\ and \\.

Example:

host=localhost port=5432 dbname=mydb connect_timeout=10

The recognized parameter key words are listed in Section 34.1.2.

34.1.1.2. Connection URIs

The general form for a connection URI is:

postgresql://[user[:password]@[netloc]][:port][,...][/dbname][?param1=value1&...]?

The URI scheme designator can be either postgresql:// or postgres://. Each of the URI parts is optional. The following examples illustrate valid URI syntax uses:

postgresql://
postgresql://localhost
postgresql://localhost:5433
postgresql://localhost/mydb
postgresql://user@localhost
postgresql://user:secret@localhost
postgresql://other@localhost/otherdb?
connect_timeout=10&application_name=myapp
postgresql://host1:123,host2:456/somedb?
target_session_attrs=any&application_name=myapp

Components of the hierarchical part of the URI can also be given as parameters. For example:

1 https://tools.ietf.org/html/rfc3986
Percent-encoding may be used to include symbols with special meaning in any of the URI parts, e.g., replace \(=\) with \%3D.

Any connection parameters not corresponding to key words listed in Section 34.1.2 are ignored and a warning message about them is sent to stderr.

For improved compatibility with JDBC connection URIs, instances of parameter \(ssl=true\) are translated into \(sslmode=require\).

The host part may be either host name or an IP address. To specify an IPv6 host address, enclose it in square brackets:

```
postgresql://[2001:db8::1234]/database
```

The host component is interpreted as described for the parameter host. In particular, a Unix-domain socket connection is chosen if the host part is either empty or starts with a slash, otherwise a TCP/IP connection is initiated. Note, however, that the slash is a reserved character in the hierarchical part of the URI. So, to specify a non-standard Unix-domain socket directory, either omit the host specification in the URI and specify the host as a parameter, or percent-encode the path in the host component of the URI:

```
postgresql:///dbname?host=/var/lib/postgresql
postgresql://%2Fvar%2Flib%2FpostgreSQL/dbname
```

It is possible to specify multiple host components, each with an optional port component, in a single URI. A URI of the form `postgresql://host1:port1,host2:port2,host3:port3/` is equivalent to a connection string of the form `host=host1,host2,host3 port=port1,port2,port3`. Each host will be tried in turn until a connection is successfully established.

### 34.1.1.3. Specifying Multiple Hosts

It is possible to specify multiple hosts to connect to, so that they are tried in the given order. In the Keyword/Value format, the `host`, `hostaddr`, and `port` options accept a comma-separated list of values. The same number of elements must be given in each option that is specified, such that e.g., the first `hostaddr` corresponds to the first host name, the second `hostaddr` corresponds to the second host name, and so forth. As an exception, if only one `port` is specified, it applies to all the hosts.

In the connection URI format, you can list multiple `host:port` pairs separated by commas, in the `host` component of the URI.

In either format, a single host name can translate to multiple network addresses. A common example of this is a host that has both an IPv4 and an IPv6 address.

When multiple hosts are specified, or when a single host name is translated to multiple addresses, all the hosts and addresses will be tried in order, until one succeeds. If none of the hosts can be reached, the connection fails. If a connection is established successfully, but authentication fails, the remaining hosts in the list are not tried.

If a password file is used, you can have different passwords for different hosts. All the other connection options are the same for every host in the list; it is not possible to e.g., specify different usernames for different hosts.

### 34.1.2. Parameter Key Words

The currently recognized parameter key words are:
**host**

Name of host to connect to. If a host name begins with a slash, it specifies Unix-domain communication rather than TCP/IP communication; the value is the name of the directory in which the socket file is stored. The default behavior when host is not specified, or is empty, is to connect to a Unix-domain socket in /tmp (or whatever socket directory was specified when PostgreSQL was built). On machines without Unix-domain sockets, the default is to connect to localhost.

A comma-separated list of host names is also accepted, in which case each host name in the list is tried in order; an empty item in the list selects the default behavior as explained above. See Section 34.1.1.3 for details.

**hostaddr**

Numeric IP address of host to connect to. This should be in the standard IPv4 address format, e.g., 172.28.40.9. If your machine supports IPv6, you can also use those addresses. TCP/IP communication is always used when a nonempty string is specified for this parameter.

Using hostaddr instead of host allows the application to avoid a host name look-up, which might be important in applications with time constraints. However, a host name is required for GSSAPI or SSPI authentication methods, as well as for verify-full SSL certificate verification. The following rules are used:

- If host is specified without hostaddr, a host name lookup occurs. (When using PQconnectPoll, the lookup occurs when PQconnectPoll first considers this host name, and it may cause PQconnectPoll to block for a significant amount of time.)
- If hostaddr is specified without host, the value for hostaddr gives the server network address. The connection attempt will fail if the authentication method requires a host name.
- If both host and hostaddr are specified, the value for hostaddr gives the server network address. The value for host is ignored unless the authentication method requires it, in which case it will be used as the host name.

Note that authentication is likely to fail if host is not the name of the server at network address hostaddr. Also, when both host and hostaddr are specified, host is used to identify the connection in a password file (see Section 34.15).

A comma-separated list of hostaddr values is also accepted, in which case each host in the list is tried in order. An empty item in the list causes the corresponding host name to be used, or the default host name if that is empty as well. See Section 34.1.1.3 for details.

Without either a host name or host address, libpq will connect using a local Unix-domain socket; or on machines without Unix-domain sockets, it will attempt to connect to localhost.

**port**

Port number to connect to at the server host, or socket file name extension for Unix-domain connections. If multiple hosts were given in the host or hostaddr parameters, this parameter may specify a comma-separated list of ports of the same length as the host list, or it may specify a single port number to be used for all hosts. An empty string, or an empty item in a comma-separated list, specifies the default port number established when PostgreSQL was built.

**dbname**

The database name. Defaults to be the same as the user name. In certain contexts, the value is checked for extended formats; see Section 34.1.1 for more details on those.

**user**

PostgreSQL user name to connect as. Defaults to be the same as the operating system name of the user running the application.
password

Password to be used if the server demands password authentication.

passfile

Specifies the name of the file used to store passwords (see Section 34.15). Defaults to ~/.pgpass, or %APPDATA%\postgresql\pgpass.conf on Microsoft Windows. (No error is reported if this file does not exist.)

connect_timeout

Maximum wait for connection, in seconds (write as a decimal integer, e.g. 10). Zero, negative, or not specified means wait indefinitely. The minimum allowed timeout is 2 seconds, therefore a value of 1 is interpreted as 2. This timeout applies separately to each host name or IP address. For example, if you specify two hosts and connect_timeout is 5, each host will time out if no connection is made within 5 seconds, so the total time spent waiting for a connection might be up to 10 seconds.

client_encoding

This sets the client_encoding configuration parameter for this connection. In addition to the values accepted by the corresponding server option, you can use auto to determine the right encoding from the current locale in the client (LC_CTYPE environment variable on Unix systems).

options

Specifies command-line options to send to the server at connection start. For example, setting this to -c geqo=off sets the session's value of the geqo parameter to off. Spaces within this string are considered to separate command-line arguments, unless escaped with a backslash (\); write \ to represent a literal backslash. For a detailed discussion of the available options, consult Chapter 19.

application_name

Specifies a value for the application_name configuration parameter.

fallback_application_name

Specifies a fallback value for the application_name configuration parameter. This value will be used if no value has been given for application_name via a connection parameter or the PGAPPNAME environment variable. Specifying a fallback name is useful in generic utility programs that wish to set a default application name but allow it to be overridden by the user.

keepalives

Controls whether client-side TCP keepalives are used. The default value is 1, meaning on, but you can change this to 0, meaning off, if keepalives are not wanted. This parameter is ignored for connections made via a Unix-domain socket.

keepalives_idle

Controls the number of seconds of inactivity after which TCP should send a keepalive message to the server. A value of zero uses the system default. This parameter is ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPIDLE or an equivalent socket option is available, and on Windows; on other systems, it has no effect.

keepalives_interval

Controls the number of seconds after which a TCP keepalive message that is not acknowledged by the server should be retransmitted. A value of zero uses the system default. This parameter is
ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPINTVL or an equivalent socket option is available, and on Windows; on other systems, it has no effect.

**keepalives_count**

Controls the number of TCP keepalives that can be lost before the client’s connection to the server is considered dead. A value of zero uses the system default. This parameter is ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPCNT or an equivalent socket option is available; on other systems, it has no effect.

**tty**

Ignored (formerly, this specified where to send server debug output).

**replication**

This option determines whether the connection should use the replication protocol instead of the normal protocol. This is what PostgreSQL replication connections as well as tools such as pg_basebackup use internally, but it can also be used by third-party applications. For a description of the replication protocol, consult Section 53.4.

The following values, which are case-insensitive, are supported:

- true, on, yes, 1
  - The connection goes into physical replication mode.
- database
  - The connection goes into logical replication mode, connecting to the database specified in the dbname parameter.
- false, off, no, 0
  - The connection is a regular one, which is the default behavior.

In physical or logical replication mode, only the simple query protocol can be used.

**sslmode**

This option determines whether or with what priority a secure SSL TCP/IP connection will be negotiated with the server. There are six modes:

- disable
  - only try a non-SSL connection
- allow
  - first try a non-SSL connection; if that fails, try an SSL connection
- prefer (default)
  - first try an SSL connection; if that fails, try a non-SSL connection
- require
  - only try an SSL connection. If a root CA file is present, verify the certificate in the same way as if verify-ca was specified
verify-ca

only try an SSL connection, and verify that the server certificate is issued by a trusted certificate authority (CA)

verify-full

only try an SSL connection, verify that the server certificate is issued by a trusted CA and that the requested server host name matches that in the certificate

See Section 34.18 for a detailed description of how these options work.

sslmode is ignored for Unix domain socket communication. If PostgreSQL is compiled without SSL support, using options require, verify-ca, or verify-full will cause an error, while options allow and prefer will be accepted but libpq will not actually attempt an SSL connection.

requiressl

This option is deprecated in favor of the sslmode setting.

If set to 1, an SSL connection to the server is required (this is equivalent to sslmode require). libpq will then refuse to connect if the server does not accept an SSL connection. If set to 0 (default), libpq will negotiate the connection type with the server (equivalent to sslmode prefer). This option is only available if PostgreSQL is compiled with SSL support.

sslcompression

If set to 1, data sent over SSL connections will be compressed. If set to 0, compression will be disabled. The default is 0. This parameter is ignored if a connection without SSL is made.

SSL compression is nowadays considered insecure and its use is no longer recommended. OpenSSL 1.1.0 disables compression by default, and many operating system distributions disable it in prior versions as well, so setting this parameter to on will not have any effect if the server does not accept compression. On the other hand, OpenSSL before 1.0.0 does not support disabling compression, so this parameter is ignored with those versions, and whether compression is used depends on the server.

If security is not a primary concern, compression can improve throughput if the network is the bottleneck. Disabling compression can improve response time and throughput if CPU performance is the limiting factor.

sslcert

This parameter specifies the file name of the client SSL certificate, replacing the default ~/.postgresql/postgresql.crt. This parameter is ignored if an SSL connection is not made.

sslkey

This parameter specifies the location for the secret key used for the client certificate. It can either specify a file name that will be used instead of the default ~/.postgresql/postgresql.key, or it can specify a key obtained from an external "engine" (engines are OpenSSL loadable modules). An external engine specification should consist of a colon-separated engine name and an engine-specific key identifier. This parameter is ignored if an SSL connection is not made.

sslrootcert

This parameter specifies the name of a file containing SSL certificate authority (CA) certificate(s). If the file exists, the server's certificate will be verified to be signed by one of these authorities. The default is ~/.postgresql/root.crt.
sslcrl

This parameter specifies the file name of the SSL certificate revocation list (CRL). Certificates listed in this file, if it exists, will be rejected while attempting to authenticate the server’s certificate. The default is ~/.postgresql/root.crl.

requirepeer

This parameter specifies the operating-system user name of the server, for example requirepeer=postgres. When making a Unix-domain socket connection, if this parameter is set, the client checks at the beginning of the connection that the server process is running under the specified user name; if it is not, the connection is aborted with an error. This parameter can be used to provide server authentication similar to that available with SSL certificates on TCP/IP connections. (Note that if the Unix-domain socket is in /tmp or another publicly writable location, any user could start a server listening there. Use this parameter to ensure that you are connected to a server run by a trusted user.) This option is only supported on platforms for which the peer authentication method is implemented; see Section 20.9.

krbsrvname

Kerberos service name to use when authenticating with GSSAPI. This must match the service name specified in the server configuration for Kerberos authentication to succeed. (See also Section 20.6.)

gsslib

GSS library to use for GSSAPI authentication. Only used on Windows. Set to gssapi to force libpq to use the GSSAPI library for authentication instead of the default SSPI.

service

Service name to use for additional parameters. It specifies a service name in pg_service.conf that holds additional connection parameters. This allows applications to specify only a service name so connection parameters can be centrally maintained. See Section 34.16.

target_session_attrs

If this parameter is set to read-write, only a connection in which read-write transactions are accepted by default is considered acceptable. The query SHOW transaction_read_only will be sent upon any successful connection; if it returns on, the connection will be closed. If multiple hosts were specified in the connection string, any remaining servers will be tried just as if the connection attempt had failed. The default value of this parameter, any, regards all connections as acceptable.

34.2. Connection Status Functions

These functions can be used to interrogate the status of an existing database connection object.

Tip

libpq application programmers should be careful to maintain the PGconn abstraction. Use the accessor functions described below to get at the contents of PGconn. Reference to internal PGconn fields using libpq-int.h is not recommended because they are subject to change in the future.

The following functions return parameter values established at connection. These values are fixed for the life of the connection. If a multi-host connection string is used, the values of PQhost, PQport,
and PQpass can change if a new connection is established using the same PGconn object. Other values are fixed for the lifetime of the PGconn object.

PQdb

Returns the database name of the connection.

char *PQdb(const PGconn *conn);

PQuser

Returns the user name of the connection.

char *PQuser(const PGconn *conn);

PQpass

Returns the password of the connection.

char *PQpass(const PGconn *conn);

PQpass will return either the password specified in the connection parameters, or if there was none and the password was obtained from the password file, it will return that. In the latter case, if multiple hosts were specified in the connection parameters, it is not possible to rely on the result of PQpass until the connection is established. The status of the connection can be checked using the function PQstatus.

PQhost

Returns the server host name of the active connection. This can be a host name, an IP address, or a directory path if the connection is via Unix socket. (The path case can be distinguished because it will always be an absolute path, beginning with /.)

char *PQhost(const PGconn *conn);

If the connection parameters specified both host and hostaddr, then PQhost will return the host information. If only hostaddr was specified, then that is returned. If multiple hosts were specified in the connection parameters, PQhost returns the host actually connected to.

PQhost returns NULL if the conn argument is NULL. Otherwise, if there is an error producing the host information (perhaps if the connection has not been fully established or there was an error), it returns an empty string.

If multiple hosts were specified in the connection parameters, it is not possible to rely on the result of PQhost until the connection is established. The status of the connection can be checked using the function PQstatus.

PQport

Returns the port of the active connection.

char *PQport(const PGconn *conn);

If multiple ports were specified in the connection parameters, PQport returns the port actually connected to.
PQport returns NULL if the conn argument is NULL. Otherwise, if there is an error producing the port information (perhaps if the connection has not been fully established or there was an error), it returns an empty string.

If multiple ports were specified in the connection parameters, it is not possible to rely on the result of PQport until the connection is established. The status of the connection can be checked using the function PQstatus.

PQtty

Returns the debug TTY of the connection. (This is obsolete, since the server no longer pays attention to the TTY setting, but the function remains for backward compatibility.)

char *PQtty(const PGconn *conn);

PQoptions

Returns the command-line options passed in the connection request.

char *PQoptions(const PGconn *conn);

The following functions return status data that can change as operations are executed on the PGconn object.

PQstatus

Returns the status of the connection.

ConnStatusType PQstatus(const PGconn *conn);

The status can be one of a number of values. However, only two of these are seen outside of an asynchronous connection procedure: CONNECTION_OK and CONNECTION_BAD. A good connection to the database has the status CONNECTION_OK. A failed connection attempt is signaled by status CONNECTION_BAD. Ordinarily, an OK status will remain so until PQfinish, but a communications failure might result in the status changing to CONNECTION_BAD prematurely. In that case the application could try to recover by calling PQreset.

See the entry for PQconnectStartParams, PQconnectStart and PQconnectPoll with regards to other status codes that might be returned.

PQtransactionStatus

Returns the current in-transaction status of the server.

PGTransactionStatusType PQtransactionStatus(const PGconn *conn);

The status can be PQTRANS_IDLE (currently idle), PQTRANS_ACTIVE (a command is in progress), PQTRANS_INTRANS (idle, in a valid transaction block), or PQTRANS_INERROR (idle, in a failed transaction block). PQTRANS_UNKNOWN is reported if the connection is bad. PQTRANS_ACTIVE is reported only when a query has been sent to the server and not yet completed.

PQparameterStatus

 Looks up a current parameter setting of the server.
const char *PQparameterStatus(const PGconn *conn, const char *paramName);

Certain parameter values are reported by the server automatically at connection startup or whenever their values change. PQparameterStatus can be used to interrogate these settings. It returns the current value of a parameter if known, or NULL if the parameter is not known.

Parameters reported as of the current release include server_version, server_encoding, client_encoding, application_name, is_superuser, session_authorization, DateStyle, IntervalStyle, TimeZone, integer_datetimes, and standard_conforming_strings. (server_encoding, TimeZone, and integer_datetimes were not reported by releases before 8.0; standard_conforming_strings was not reported by releases before 8.1; IntervalStyle was not reported by releases before 8.4; application_name was not reported by releases before 9.0.) Note that server_version, server_encoding and integer_datetimes cannot change after startup.

Pre-3.0-protocol servers do not report parameter settings, but libpq includes logic to obtain values for server_version and client_encoding anyway. Applications are encouraged to use PQparameterStatus rather than ad hoc code to determine these values. (Beware however that on a pre-3.0 connection, changing client_encoding via SET after connection startup will not be reflected by PQparameterStatus.) For server_version, see also PQserverVersion, which returns the information in a numeric form that is much easier to compare against.

If no value for standard_conforming_strings is reported, applications can assume it is off, that is, backslashes are treated as escapes in string literals. Also, the presence of this parameter can be taken as an indication that the escape string syntax (E'...') is accepted.

Although the returned pointer is declared const, it in fact points to mutable storage associated with the PGconn structure. It is unwise to assume the pointer will remain valid across queries.

PQprotocolVersion

Interrogates the frontend/backend protocol being used.

int PQprotocolVersion(const PGconn *conn);

Applications might wish to use this function to determine whether certain features are supported. Currently, the possible values are 2 (2.0 protocol), 3 (3.0 protocol), or zero (connection bad). The protocol version will not change after connection startup is complete, but it could theoretically change during a connection reset. The 3.0 protocol will normally be used when communicating with PostgreSQL 7.4 or later servers; pre-7.4 servers support only protocol 2.0. (Protocol 1.0 is obsolete and not supported by libpq.)

PQserverVersion

Returns an integer representing the server version.

int PQserverVersion(const PGconn *conn);

Applications might use this function to determine the version of the database server they are connected to. The result is formed by multiplying the server's major version number by 10000 and adding the minor version number. For example, version 10.1 will be returned as 100001, and version 11.0 will be returned as 110000. Zero is returned if the connection is bad.

Prior to major version 10, PostgreSQL used three-part version numbers in which the first two parts together represented the major version. For those versions, PQserverVersion uses two
digits for each part; for example version 9.1.5 will be returned as 90105, and version 9.2.0 will be returned as 90200.

Therefore, for purposes of determining feature compatibility, applications should divide the result of PQserverVersion by 100 not 10000 to determine a logical major version number. In all release series, only the last two digits differ between minor releases (bug-fix releases).

PQerrorMessage

Returns the error message most recently generated by an operation on the connection.

char *PQerrorMessage(const PGconn *conn);

Nearly all libpq functions will set a message for PQerrorMessage if they fail. Note that by libpq convention, a nonempty PQerrorMessage result can consist of multiple lines, and will include a trailing newline. The caller should not free the result directly. It will be freed when the associated PGconn handle is passed to PQfinish. The result string should not be expected to remain the same across operations on the PGconn structure.

PQsocket

Obtains the file descriptor number of the connection socket to the server. A valid descriptor will be greater than or equal to 0; a result of -1 indicates that no server connection is currently open. (This will not change during normal operation, but could change during connection setup or reset.)

int PQsocket(const PGconn *conn);

PQbackendPID

Returns the process ID (PID) of the backend process handling this connection.

int PQbackendPID(const PGconn *conn);

The backend PID is useful for debugging purposes and for comparison to NOTIFY messages (which include the PID of the notifying backend process). Note that the PID belongs to a process executing on the database server host, not the local host!

PQconnectionNeedsPassword

Returns true (1) if the connection authentication method required a password, but none was available. Returns false (0) if not.

int PQconnectionNeedsPassword(const PGconn *conn);

This function can be applied after a failed connection attempt to decide whether to prompt the user for a password.

PQconnectionUsedPassword

Returns true (1) if the connection authentication method used a password. Returns false (0) if not.

int PQconnectionUsedPassword(const PGconn *conn);

This function can be applied after either a failed or successful connection attempt to detect whether the server demanded a password.
The following functions return information related to SSL. This information usually doesn't change after a connection is established.

**PQsslInUse**

Returns true (1) if the connection uses SSL, false (0) if not.

```c
int PQsslInUse(const PGconn *conn);
```

**PQsslAttribute**

Returns SSL-related information about the connection.

```c
const char *PQsslAttribute(const PGconn *conn, const char *attribute_name);
```

The list of available attributes varies depending on the SSL library being used, and the type of connection. If an attribute is not available, returns NULL.

The following attributes are commonly available:

- **library**
  
  Name of the SSL implementation in use. (Currently, only "OpenSSL" is implemented)

- **protocol**
  
  SSL/TLS version in use. Common values are "TLSv1", "TLSv1.1" and "TLSv1.2", but an implementation may return other strings if some other protocol is used.

- **key_bits**
  
  Number of key bits used by the encryption algorithm.

- **cipher**
  
  A short name of the ciphersuite used, e.g. "DHE-RSA-DES-CBC3-SHA". The names are specific to each SSL implementation.

- **compression**
  
  If SSL compression is in use, returns the name of the compression algorithm, or "on" if compression is used but the algorithm is not known. If compression is not in use, returns "off".

**PQsslAttributeNames**

Return an array of SSL attribute names available. The array is terminated by a NULL pointer.

```c
const char *const * PQsslAttributeNames(const PGconn *conn);
```

**PQsslStruct**

Return a pointer to an SSL-implementation-specific object describing the connection.

```c
void *PQsslStruct(const PGconn *conn, const char *struct_name);
```
The struct(s) available depend on the SSL implementation in use. For OpenSSL, there is one struct, available under the name "OpenSSL", and it returns a pointer to the OpenSSL SSL struct. To use this function, code along the following lines could be used:

```c
#include <libpq-fe.h>
#include <openssl/ssl.h>
...
SSL *ssl;

dbconn = PQconnectdb(...);
...
ssl = PQsslStruct(dbconn, "OpenSSL");
if (ssl)
{
    /* use OpenSSL functions to access ssl */
}
```
This structure can be used to verify encryption levels, check server certificates, and more. Refer to the OpenSSL documentation for information about this structure.

**PQgetssl**

Returns the SSL structure used in the connection, or null if SSL is not in use.

```c
void *PQgetssl(const PGconn *conn);
```
This function is equivalent to `PQsslStruct(conn, "OpenSSL")`. It should not be used in new applications, because the returned struct is specific to OpenSSL and will not be available if another SSL implementation is used. To check if a connection uses SSL, call `PQsslInUse` instead, and for more details about the connection, use `PQsslAttribute`.

### 34.3. Command Execution Functions

Once a connection to a database server has been successfully established, the functions described here are used to perform SQL queries and commands.

#### 34.3.1. Main Functions

**PQexec**

Submits a command to the server and waits for the result.

```c
PGresult *PQexec(PGconn *conn, const char *command);
```
Returns a `PGresult` pointer or possibly a null pointer. A non-null pointer will generally be returned except in out-of-memory conditions or serious errors such as inability to send the command to the server. The `PQresultStatus` function should be called to check the return value for any errors (including the value of a null pointer, in which case it will return `PGRES_FATAL_ERROR`). Use `PQerrorMessage` to get more information about such errors.

The command string can include multiple SQL commands (separated by semicolons). Multiple queries sent in a single `PQexec` call are processed in a single transaction, unless there are explicit `BEGIN/...`
COMMIT commands included in the query string to divide it into multiple transactions. (See Section 53.2.2.1 for more details about how the server handles multi-query strings.) Note however that the returned PGresult structure describes only the result of the last command executed from the string. Should one of the commands fail, processing of the string stops with it and the returned PGresult describes the error condition.

PQexecParams

Submits a command to the server and waits for the result, with the ability to pass parameters separately from the SQL command text.

PGresult *PQexecParams(PGconn *conn,
            const char *command,
            int nParams,
            const Oid *paramTypes,
            const char * const *paramValues,
            const int *paramLengths,
            const int *paramFormats,
            int resultFormat);

PQexecParams is like PQexec, but offers additional functionality: parameter values can be specified separately from the command string proper, and query results can be requested in either text or binary format. PQexecParams is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

The function arguments are:

conn

The connection object to send the command through.

command

The SQL command string to be executed. If parameters are used, they are referred to in the command string as $1, $2, etc.

nParams

The number of parameters supplied; it is the length of the arrays paramTypes[], paramValues[], paramLengths[], and paramFormats[]. (The array pointers can be NULL when nParams is zero.)

paramTypes[]

Specifies, by OID, the data types to be assigned to the parameter symbols. If paramTypes is NULL, or any particular element in the array is zero, the server infers a data type for the parameter symbol in the same way it would do for an untyped literal string.

paramValues[]

Specifies the actual values of the parameters. A null pointer in this array means the corresponding parameter is null; otherwise the pointer points to a zero-terminated text string (for text format) or binary data in the format expected by the server (for binary format).

paramLengths[]

Specifies the actual data lengths of binary-format parameters. It is ignored for null parameters and text-format parameters. The array pointer can be null when there are no binary parameters.
**paramFormats[]**

Specifies whether parameters are text (put a zero in the array entry for the corresponding parameter) or binary (put a one in the array entry for the corresponding parameter). If the array pointer is null then all parameters are presumed to be text strings.

Values passed in binary format require knowledge of the internal representation expected by the backend. For example, integers must be passed in network byte order. Passing numeric values requires knowledge of the server storage format, as implemented in `src/backend/utils/adt/numeric.c::numeric_send()` and `src/backend/utils/adt/numeric.c::numeric_recv()`.

**resultFormat**

Specify zero to obtain results in text format, or one to obtain results in binary format. (There is not currently a provision to obtain different result columns in different formats, although that is possible in the underlying protocol.)

The primary advantage of `PQexecParams` over `PQexec` is that parameter values can be separated from the command string, thus avoiding the need for tedious and error-prone quoting and escaping.

Unlike `PQexec`, `PQexecParams` allows at most one SQL command in the given string. (There can be semicolons in it, but not more than one nonempty command.) This is a limitation of the underlying protocol, but has some usefulness as an extra defense against SQL-injection attacks.

### Tip

Specifying parameter types via OIDs is tedious, particularly if you prefer not to hardwire particular OID values into your program. However, you can avoid doing so even in cases where the server by itself cannot determine the type of the parameter, or chooses a different type than you want. In the SQL command text, attach an explicit cast to the parameter symbol to show what data type you will send. For example:

```sql
SELECT * FROM mytable WHERE x = $1::bigint;
```

This forces parameter $1 to be treated as bigint, whereas by default it would be assigned the same type as x. Forcing the parameter type decision, either this way or by specifying a numeric type OID, is strongly recommended when sending parameter values in binary format, because binary format has less redundancy than text format and so there is less chance that the server will detect a type mismatch mistake for you.

**PQprepare**

Submits a request to create a prepared statement with the given parameters, and waits for completion.

```c
PGresult *PQprepare(PGconn *conn,
            const char *stmtName,
            const char *query,
            int nParams,
            const Oid *paramTypes);
```

`PQprepare` creates a prepared statement for later execution with `PQexecPrepared`. This feature allows commands to be executed repeatedly without being parsed and planned each time; see PREPARE for details. `PQprepare` is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.
The function creates a prepared statement named \texttt{stmtName} from the \texttt{query} string, which must contain a single SQL command. \texttt{stmtName} can be "" to create an unnamed statement, in which case any pre-existing unnamed statement is automatically replaced; otherwise it is an error if the statement name is already defined in the current session. If any parameters are used, they are referred to in the query as $1, \$2, \ldots, nParams$ is the number of parameters for which types are pre-specified in the array \texttt{paramTypes[]} \texttt{(The array pointer can be NULL when nParams is zero.)} \texttt{paramTypes[]} specifies, by OID, the data types to be assigned to the parameter symbols. If \texttt{paramTypes} is NULL, or any particular element in the array is zero, the server assigns a data type to the parameter symbol in the same way it would do for an untyped literal string. Also, the query can use parameter symbols with numbers higher than \texttt{nParams}; data types will be inferred for these symbols as well. (See \texttt{PQdescribePrepared} for a means to find out what data types were inferred.)

As with \texttt{PQexec}, the result is normally a \texttt{PGresult} object whose contents indicate server-side success or failure. A null result indicates out-of-memory or inability to send the command at all. Use \texttt{PQerrorMessage} to get more information about such errors.

Prepared statements for use with \texttt{PQexecPrepared} can also be created by executing SQL PREPARE statements. Also, although there is no libpq function for deleting a prepared statement, the SQL DEALLOCATE statement can be used for that purpose.

\texttt{PQexecPrepared}

Sends a request to execute a prepared statement with given parameters, and waits for the result.

\texttt{PGresult *}\texttt{PQexecPrepared(PGconn *conn, const char *stmtName, int nParams, const char * const *paramValues, const int *paramLengths, const int *paramFormats, int resultFormat);}

\texttt{PQexecPrepared} is like \texttt{PQexecParams}, but the command to be executed is specified by naming a previously-prepared statement, instead of giving a query string. This feature allows commands that will be used repeatedly to be parsed and planned just once, rather than each time they are executed. The statement must have been prepared previously in the current session. \texttt{PQexecPrepared} is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

The parameters are identical to \texttt{PQexecParams}, except that the name of a prepared statement is given instead of a query string, and the \texttt{paramTypes[]} parameter is not present (it is not needed since the prepared statement's parameter types were determined when it was created).

\texttt{PQdescribePrepared}

Submits a request to obtain information about the specified prepared statement, and waits for completion.

\texttt{PGresult *}\texttt{PQdescribePrepared(PGconn *conn, const char *stmtName);}
MAND_OK is returned. The functions PQnparams and PQparamtype can be applied to this PGresult to obtain information about the parameters of the prepared statement, and the functions PQnfields, PQfname, PQfctype, etc provide information about the result columns (if any) of the statement.

PQdescribePortal

Submits a request to obtain information about the specified portal, and waits for completion.

PGresult *PQdescribePortal(PGconn *conn, const char *portalName);

PQdescribePortal allows an application to obtain information about a previously created portal. (libpq does not provide any direct access to portals, but you can use this function to inspect the properties of a cursor created with a DECLARE CURSOR SQL command.) PQdescribePortal is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

portalName can be "" or NULL to reference the unnamed portal, otherwise it must be the name of an existing portal. On success, a PGresult with status PGRES_COMMAND_OK is returned. The functions PQnfields, PQfname, PQfctype, etc can be applied to the PGresult to obtain information about the result columns (if any) of the portal.

The PGresult structure encapsulates the result returned by the server. libpq application programmers should be careful to maintain the PGresult abstraction. Use the accessor functions below to get at the contents of PGresult. Avoid directly referencing the fields of the PGresult structure because they are subject to change in the future.

PQresultStatus

Returns the result status of the command.

ExecStatusType PQresultStatus(const PGresult *res);

PQresultStatus can return one of the following values:

PGRES_EMPTY_QUERY

The string sent to the server was empty.

PGRES_COMMAND_OK

Successful completion of a command returning no data.

PGRES_TUPLES_OK

Successful completion of a command returning data (such as a SELECT or SHOW).

PGRES_COPY_OUT

Copy Out (from server) data transfer started.

PGRES_COPY_IN

Copy In (to server) data transfer started.

PGRES_BAD_RESPONSE

The server's response was not understood.
libpq - C Library

PGRES_NONFATAL_ERROR

A nonfatal error (a notice or warning) occurred.

PGRES_FATAL_ERROR

A fatal error occurred.

PGRES_COPY_BOTH

Copy In/Out (to and from server) data transfer started. This feature is currently used only for streaming replication, so this status should not occur in ordinary applications.

PGRES_SINGLE_TUPLE

The PGresult contains a single result tuple from the current command. This status occurs only when single-row mode has been selected for the query (see Section 34.5).

If the result status is PGRES_TUPLES_OK or PGRES_SINGLE_TUPLE, then the functions described below can be used to retrieve the rows returned by the query. Note that a SELECT command that happens to retrieve zero rows still shows PGRES_TUPLES_OK. PGRES_COMMAND_OK is for commands that can never return rows (INSERT or UPDATE without a RETURNING clause, etc.). A response of PGRES_EMPTY_QUERY might indicate a bug in the client software.

A result of status PGRES_NONFATAL_ERROR will never be returned directly by PQexec or other query execution functions; results of this kind are instead passed to the notice processor (see Section 34.12).

PQresStatus

Converts the enumerated type returned by PQresultStatus into a string constant describing the status code. The caller should not free the result.

char *PQresStatus(ExecStatusType status);

PQresultErrorMessage

Returns the error message associated with the command, or an empty string if there was no error.

char *PQresultErrorMessage(const PGresult *res);

If there was an error, the returned string will include a trailing newline. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to PQclear.

Immediately following a PQexec or PQgetResult call, PQerrorMessage (on the connection) will return the same string as PQresultErrorMessage (on the result). However, a PGresult will retain its error message until destroyed, whereas the connection’s error message will change when subsequent operations are done. Use PQresultErrorMessage when you want to know the status associated with a particular PGresult; use PQerrorMessage when you want to know the status from the latest operation on the connection.

PQresultVerboseErrorMessage

Returns a reformatted version of the error message associated with a PGresult object.

char *PQresultVerboseErrorMessage(const PGresult *res,
In some situations a client might wish to obtain a more detailed version of a previously-reported error. `PQresultVerboseErrorMessage` addresses this need by computing the message that would have been produced by `PQresultErrorMessage` if the specified verbosity settings had been in effect for the connection when the given `PGresult` was generated. If the `PGresult` is not an error result, “`PGresult is not an error result`” is reported instead. The returned string includes a trailing newline.

Unlike most other functions for extracting data from a `PGresult`, the result of this function is a freshly allocated string. The caller must free it using `PQfreemem()` when the string is no longer needed.

A NULL return is possible if there is insufficient memory.

**PQresultErrorField**

Returns an individual field of an error report.

```
char *PQresultErrorField(const PGresult *res, int fieldcode);
```

`fieldcode` is an error field identifier; see the symbols listed below. NULL is returned if the `PGresult` is not an error or warning result, or does not include the specified field. Field values will normally not include a trailing newline. The caller should not free the result directly. It will be freed when the associated `PGresult` handle is passed to `PQclear`.

The following field codes are available:

**PG_DIAG_SEVERITY**

The severity; the field contents are `ERROR`, `FATAL`, or `PANIC` (in an error message), or `WARNING`, `NOTICE`, `DEBUG`, `INFO`, or `LOG` (in a notice message), or a localized translation of one of these. Always present.

**PG_DIAG_SEVERITY_NONLOCALIZED**

The severity; the field contents are `ERROR`, `FATAL`, or `PANIC` (in an error message), or `WARNING`, `NOTICE`, `DEBUG`, `INFO`, or `LOG` (in a notice message). This is identical to the `PG_DIAG_SEVERITY` field except that the contents are never localized. This is present only in reports generated by PostgreSQL versions 9.6 and later.

**PG_DIAG_SQLSTATE**

The SQLSTATE code for the error. The SQLSTATE code identifies the type of error that has occurred; it can be used by front-end applications to perform specific operations (such as error handling) in response to a particular database error. For a list of the possible SQLSTATE codes, see Appendix A. This field is not localizable, and is always present.

**PG_DIAG_MESSAGE_PRIMARY**

The primary human-readable error message (typically one line). Always present.

**PG_DIAG_MESSAGE_DETAIL**

Detail: an optional secondary error message carrying more detail about the problem. Might run to multiple lines.
PG_DIAG_MESSAGE_HINT

Hint: an optional suggestion what to do about the problem. This is intended to differ from
detail in that it offers advice (potentially inappropriate) rather than hard facts. Might run to
multiple lines.

PG_DIAG_STATEMENT_POSITION

A string containing a decimal integer indicating an error cursor position as an index into
the original statement string. The first character has index 1, and positions are measured in
characters not bytes.

PG_DIAG_INTERNAL_POSITION

This is defined the same as the PG_DIAG_STATEMENT_POSITION field, but it is used
when the cursor position refers to an internally generated command rather than the one sub-
mitted by the client. The PG_DIAG_INTERNAL_QUERY field will always appear when this
field appears.

PG_DIAG_INTERNAL_QUERY

The text of a failed internally-generated command. This could be, for example, a SQL query
issued by a PL/pgSQL function.

PG_DIAG_CONTEXT

An indication of the context in which the error occurred. Presently this includes a call stack
traceback of active procedural language functions and internally-generated queries. The trace
is one entry per line, most recent first.

PG_DIAG_SCHEMA_NAME

If the error was associated with a specific database object, the name of the schema containing
that object, if any.

PG_DIAG_TABLE_NAME

If the error was associated with a specific table, the name of the table. (Refer to the schema
name field for the name of the table's schema.)

PG_DIAG_COLUMN_NAME

If the error was associated with a specific table column, the name of the column. (Refer to
the schema and table name fields to identify the table.)

PG_DIAG_DATATYPE_NAME

If the error was associated with a specific data type, the name of the data type. (Refer to the
schema name field for the name of the data type's schema.)

PG_DIAG_CONSTRAINT_NAME

If the error was associated with a specific constraint, the name of the constraint. Refer to
fields listed above for the associated table or domain. (For this purpose, indexes are treated
as constraints, even if they weren't created with constraint syntax.)

PG_DIAG_SOURCE_FILE

The file name of the source-code location where the error was reported.

PG_DIAG_SOURCE_LINE

The line number of the source-code location where the error was reported.
The name of the source-code function reporting the error.

**Note**

The fields for schema name, table name, column name, data type name, and constraint name are supplied only for a limited number of error types; see Appendix A. Do not assume that the presence of any of these fields guarantees the presence of another field. Core error sources observe the interrelationships noted above, but user-defined functions may use these fields in other ways. In the same vein, do not assume that these fields denote contemporary objects in the current database.

The client is responsible for formatting displayed information to meet its needs; in particular it should break long lines as needed. Newline characters appearing in the error message fields should be treated as paragraph breaks, not line breaks.

Errors generated internally by libpq will have severity and primary message, but typically no other fields. Errors returned by a pre-3.0-protocol server will include severity and primary message, and sometimes a detail message, but no other fields.

Note that error fields are only available from `PGresult` objects, not `PGconn` objects; there is no `PQerrorField` function.

**PQclear**

Frees the storage associated with a `PGresult`. Every command result should be freed via `PQclear` when it is no longer needed.

```c
void PQclear(PGresult *res);
```

You can keep a `PGresult` object around for as long as you need it; it does not go away when you issue a new command, nor even if you close the connection. To get rid of it, you must call `PQclear`. Failure to do this will result in memory leaks in your application.

### 34.3.2. Retrieving Query Result Information

These functions are used to extract information from a `PGresult` object that represents a successful query result (that is, one that has status `PGRES_TUPLES_OK` or `PGRES_SINGLE_TUPLE`). They can also be used to extract information from a successful Describe operation: a Describe's result has all the same column information that actual execution of the query would provide, but it has zero rows. For objects with other status values, these functions will act as though the result has zero rows and zero columns.

**PQntuples**

Returns the number of rows (tuples) in the query result. (Note that `PGresult` objects are limited to no more than `INT_MAX` rows, so an `int` result is sufficient.)

```c
int PQntuples(const PGresult *res);
```

**PQnfields**

Returns the number of columns (fields) in each row of the query result.
int PQnfields(const PGresult *res);

PQfname

Returns the column name associated with the given column number. Column numbers start at 0. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to PQclear.

char *PQfname(const PGresult *res,
int column_number);

NULL is returned if the column number is out of range.

PQfnumber

Returns the column number associated with the given column name.

int PQfnumber(const PGresult *res,
const char *column_name);

-1 is returned if the given name does not match any column.

The given name is treated like an identifier in an SQL command, that is, it is downcased unless double-quoted. For example, given a query result generated from the SQL command:

SELECT 1 AS FOO, 2 AS "BAR";

we would have the results:

PQfname(res, 0)       foo
PQfname(res, 1)       BAR
PQfnumber(res, "FOO") 0
PQfnumber(res, "foo") 0
PQfnumber(res, "BAR") -1
PQfnumber(res, "\"BAR\"") 1

PQftable

Returns the OID of the table from which the given column was fetched. Column numbers start at 0.

Oid PQftable(const PGresult *res,
int column_number);

InvalidOid is returned if the column number is out of range, or if the specified column is not a simple reference to a table column, or when using pre-3.0 protocol. You can query the system table pg_class to determine exactly which table is referenced.

The type Oid and the constant InvalidOid will be defined when you include the libpq header file. They will both be some integer type.

PQftablecol

Returns the column number (within its table) of the column making up the specified query result column. Query-result column numbers start at 0, but table columns have nonzero numbers.
libpq - C Library

```
int PQftablecol(const PGresult *res,
                 int column_number);
```

Zero is returned if the column number is out of range, or if the specified column is not a simple reference to a table column, or when using pre-3.0 protocol.

**PQfformat**

Returns the format code indicating the format of the given column. Column numbers start at 0.

```
int PQfformat(const PGresult *res,
              int column_number);
```

Format code zero indicates textual data representation, while format code one indicates binary representation. (Other codes are reserved for future definition.)

**PQftype**

Returns the data type associated with the given column number. The integer returned is the internal OID number of the type. Column numbers start at 0.

```
Oid PQftype(const PGresult *res,
            int column_number);
```

You can query the system table `pg_type` to obtain the names and properties of the various data types. The OIDs of the built-in data types are defined in the file `src/include/catalog/pg_type_d.h` in the source tree.

**PQfmod**

Returns the type modifier of the column associated with the given column number. Column numbers start at 0.

```
int PQfmod(const PGresult *res,
           int column_number);
```

The interpretation of modifier values is type-specific; they typically indicate precision or size limits. The value -1 is used to indicate “no information available”. Most data types do not use modifiers, in which case the value is always -1.

**PQfsize**

Returns the size in bytes of the column associated with the given column number. Column numbers start at 0.

```
int PQfsize(const PGresult *res,
            int column_number);
```

`PQfsize` returns the space allocated for this column in a database row, in other words the size of the server's internal representation of the data type. (Accordingly, it is not really very useful to clients.) A negative value indicates the data type is variable-length.

**PQbinaryTuples**

Returns 1 if the `PGresult` contains binary data and 0 if it contains text data.
int PQbinaryTuples(const PGresult *res);

This function is deprecated (except for its use in connection with \texttt{COPY}), because it is possible for a single PGresult to contain text data in some columns and binary data in others. \texttt{PQformat} is preferred. \texttt{PQbinaryTuples} returns 1 only if all columns of the result are binary (format 1).

**PQgetvalue**

Returns a single field value of one row of a PGresult. Row and column numbers start at 0. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to \texttt{PQclear}.

```c
char *PQgetvalue(const PGresult *res,
                 int row_number,
                 int column_number);
```

For data in text format, the value returned by \texttt{PQgetvalue} is a null-terminated character string representation of the field value. For data in binary format, the value is in the binary representation determined by the data type's \texttt{typsend} and \texttt{typreceive} functions. (The value is actually followed by a zero byte in this case too, but that is not ordinarily useful, since the value is likely to contain embedded nulls.)

An empty string is returned if the field value is null. See \texttt{PQgetisnull} to distinguish null values from empty-string values.

The pointer returned by \texttt{PQgetvalue} points to storage that is part of the PGresult structure. One should not modify the data it points to, and one must explicitly copy the data into other storage if it is to be used past the lifetime of the PGresult structure itself.

**PQgetisnull**

Tests a field for a null value. Row and column numbers start at 0.

```c
int PQgetisnull(const PGresult *res,
                 int row_number,
                 int column_number);
```

This function returns 1 if the field is null and 0 if it contains a non-null value. (Note that \texttt{PQgetvalue} will return an empty string, not a null pointer, for a null field.)

**PQgetlength**

Returns the actual length of a field value in bytes. Row and column numbers start at 0.

```c
int PQgetlength(const PGresult *res,
                 int row_number,
                 int column_number);
```

This is the actual data length for the particular data value, that is, the size of the object pointed to by \texttt{PQgetvalue}. For text data format this is the same as \texttt{strlen()}. For binary format this is essential information. Note that one should not rely on \texttt{PQfsize} to obtain the actual data length.

**PQnparams**

Returns the number of parameters of a prepared statement.
int PQnparams(const PGresult *res);

This function is only useful when inspecting the result of PQdescribePrepared. For other types of queries it will return zero.

PQparamtype

Returns the data type of the indicated statement parameter. Parameter numbers start at 0.

Old PQparamtype(const PGresult *res, int param_number);

This function is only useful when inspecting the result of PQdescribePrepared. For other types of queries it will return zero.

PQprint

Prints out all the rows and, optionally, the column names to the specified output stream.

void PQprint(FILE *fout, /* output stream */
              const PGresult *res,
              const PQprintOpt *po);

typedef struct
{
    pqbool  header; /* print output field headings and row count */
    pqbool  align; /* fill align the fields */
    pqbool  standard; /* old brain dead format */
    pqbool  html3; /* output HTML tables */
    pqbool  expanded; /* expand tables */
    pqbool  pager; /* use pager for output if needed */
    char   *fieldSep; /* field separator */
    char   *tableOpt; /* attributes for HTML table element */
    char   *caption; /* HTML table caption */
    char   **fieldName; /* null-terminated array of replacement field names */
} PQprintOpt;

This function was formerly used by psql to print query results, but this is no longer the case. Note that it assumes all the data is in text format.

34.3.3. Retrieving Other Result Information

These functions are used to extract other information from PGresult objects.

PQcmdStatus

Returns the command status tag from the SQL command that generated the PGresult.

char *PQcmdStatus(PGresult *res);

Commonly this is just the name of the command, but it might include additional data such as the number of rows processed. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to PQclear.

PQcmdTuples

Returns the number of rows affected by the SQL command.
This function returns a string containing the number of rows affected by the SQL statement that generated the PGresult. This function can only be used following the execution of a SELECT, CREATE TABLE AS, INSERT, UPDATE, DELETE, MOVE, FETCH, or COPY statement, or an EXECUTE of a prepared query that contains an INSERT, UPDATE, or DELETE statement. If the command that generated the PGresult was anything else, PQcmdTuples returns an empty string. The caller should not free the return value directly. It will be freed when the associated PGresult handle is passed to PQclear.

PQoidValue

Returns the OID of the inserted row, if the SQL command was an INSERT that inserted exactly one row into a table that has OIDs, or a EXECUTE of a prepared query containing a suitable INSERT statement. Otherwise, this function returns InvalidOid. This function will also return InvalidOid if the table affected by the INSERT statement does not contain OIDs.

Oid PQoidValue(const PGresult *res);

PQoidStatus

This function is deprecated in favor of PQoidValue and is not thread-safe. It returns a string with the OID of the inserted row, while PQoidValue returns the OID value.

char *PQoidStatus(const PGresult *res);

34.3.4. Escaping Strings for Inclusion in SQL Commands

PQescapeLiteral

char *PQescapeLiteral(PGconn *conn, const char *str, size_t length);

PQescapeLiteral escapes a string for use within an SQL command. This is useful when inserting data values as literal constants in SQL commands. Certain characters (such as quotes and backslashes) must be escaped to prevent them from being interpreted specially by the SQL parser. PQescapeLiteral performs this operation.

PQescapeLiteral returns an escaped version of the str parameter in memory allocated with malloc(). This memory should be freed using PQfreemem() when the result is no longer needed. A terminating zero byte is not required, and should not be counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeLiteral stops at the zero; the behavior is thus rather like strncpy.) The return string has all special characters replaced so that they can be properly processed by the PostgreSQL string literal parser. A terminating zero byte is also added. The single quotes that must surround PostgreSQL string literals are included in the result string.

On error, PQescapeLiteral returns NULL and a suitable message is stored in the conn object.

Tip

It is especially important to do proper escaping when handling strings that were received from an untrustworthy source. Otherwise there is a security risk: you are
libpq - C Library

vulnerable to “SQL injection” attacks wherein unwanted SQL commands are fed to your database.

Note that it is not necessary nor correct to do escaping when a data value is passed as a separate parameter in PQexecParams or its sibling routines.

PQescapeIdentifier

char *PQescapeIdentifier(PGconn *conn, const char *str, size_t length);

PQescapeIdentifier escapes a string for use as an SQL identifier, such as a table, column, or function name. This is useful when a user-supplied identifier might contain special characters that would otherwise not be interpreted as part of the identifier by the SQL parser, or when the identifier might contain upper case characters whose case should be preserved.

PQescapeIdentifier returns a version of the str parameter escaped as an SQL identifier in memory allocated with malloc(). This memory must be freed using PQfreemem() when the result is no longer needed. A terminating zero byte is not required, and should not be counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeIdentifier stops at the zero; the behavior is thus rather like strncpy.) The return string has all special characters replaced so that it will be properly processed as an SQL identifier. A terminating zero byte is also added. The return string will also be surrounded by double quotes.

On error, PQescapeIdentifier returns NULL and a suitable message is stored in the conn object.

Tip

As with string literals, to prevent SQL injection attacks, SQL identifiers must be escaped when they are received from an untrustworthy source.

PQescapeStringConn

size_t PQescapeStringConn(PGconn *conn, char *to, const char *from, size_t length, int *error);

PQescapeStringConn escapes string literals, much like PQescapeLiteral. Unlike PQescapeLiteral, the caller is responsible for providing an appropriately sized buffer. Furthermore, PQescapeStringConn does not generate the single quotes that must surround PostgreSQL string literals; they should be provided in the SQL command that the result is inserted into. The parameter from points to the first character of the string that is to be escaped, and the length parameter gives the number of bytes in this string. A terminating zero byte is not required, and should not be counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeStringConn stops at the zero; the behavior is thus rather like strncpy.) to shall point to a buffer that is able to hold at least one more byte than twice the value of length, otherwise the behavior is undefined. Behavior is likewise undefined if the to and from strings overlap.

If the error parameter is not NULL, then *error is set to zero on success, nonzero on error. Presently the only possible error conditions involve invalid multibyte encoding in the source string. The output string is still generated on error, but it can be expected that the server will
libpq - C Library

reject it as malformed. On error, a suitable message is stored in the conn object, whether or not error is NULL.

PQescapeStringConn returns the number of bytes written to to, not including the terminating zero byte.

PQescapeString

PQescapeString is an older, deprecated version of PQescapeStringConn.

size_t PQescapeString (char *to, const char *from, size_t length);

The only difference from PQescapeStringConn is that PQescapeString does not take PGconn or error parameters. Because of this, it cannot adjust its behavior depending on the connection properties (such as character encoding) and therefore it might give the wrong results. Also, it has no way to report error conditions.

PQescapeString can be used safely in client programs that work with only one PostgreSQL connection at a time (in this case it can find out what it needs to know "behind the scenes"). In other contexts it is a security hazard and should be avoided in favor of PQescapeStringConn.

PQescapeByteaConn

Escapes binary data for use within an SQL command with the type bytea. As with PQescapeStringConn, this is only used when inserting data directly into an SQL command string.

unsigned char *PQescapeByteaConn(PGconn *conn, const unsigned char *from, size_t from_length, size_t *to_length);

Certain byte values must be escaped when used as part of a bytea literal in an SQL statement. PQescapeByteaConn escapes bytes using either hex encoding or backslash escaping. See Section 8.4 for more information.

The from parameter points to the first byte of the string that is to be escaped, and the from_length parameter gives the number of bytes in this binary string. (A terminating zero byte is neither necessary nor counted.) The to_length parameter points to a variable that will hold the resultant escaped string length. This result string length includes the terminating zero byte of the result.

PQescapeByteaConn returns an escaped version of the from parameter binary string in memory allocated with malloc(). This memory should be freed using PQfreemem() when the result is no longer needed. The return string has all special characters replaced so that they can be properly processed by the PostgreSQL string literal parser, and the bytea input function. A terminating zero byte is also added. The single quotes that must surround PostgreSQL string literals are not part of the result string.

On error, a null pointer is returned, and a suitable error message is stored in the conn object. Currently, the only possible error is insufficient memory for the result string.

PQescapeBytea

PQescapeBytea is an older, deprecated version of PQescapeByteaConn.

unsigned char *PQescapeBytea(const unsigned char *from,
size_t from_length,
size_t *to_length);

The only difference from PQescapeByteaConn is that PQescapeBytea does not take a PGconn parameter. Because of this, PQescapeBytea can only be used safely in client programs that use a single PostgreSQL connection at a time (in this case it can find out what it needs to know “behind the scenes”). It might give the wrong results if used in programs that use multiple database connections (use PQescapeByteaConn in such cases).

PQunescapeBytea

Converts a string representation of binary data into binary data — the reverse of PQescapeBytea. This is needed when retrieving bytea data in text format, but not when retrieving it in binary format.

unsigned char *PQunescapeBytea(const unsigned char *from, size_t *to_length);

The from parameter points to a string such as might be returned by PQgetvalue when applied to a bytea column. PQunescapeBytea converts this string representation into its binary representation. It returns a pointer to a buffer allocated with malloc(), or NULL on error, and puts the size of the buffer in to_length. The result must be freed using PQfreemem when it is no longer needed.

This conversion is not exactly the inverse of PQescapeBytea, because the string is not expected to be “escaped” when received from PQgetvalue. In particular this means there is no need for string quoting considerations, and so no need for a PGconn parameter.

34.4. Asynchronous Command Processing

The PQexec function is adequate for submitting commands in normal, synchronous applications. It has a few deficiencies, however, that can be of importance to some users:

• PQexec waits for the command to be completed. The application might have other work to do (such as maintaining a user interface), in which case it won’t want to block waiting for the response.

• Since the execution of the client application is suspended while it waits for the result, it is hard for the application to decide that it would like to try to cancel the ongoing command. (It can be done from a signal handler, but not otherwise.)

• PQexec can return only one PGresult structure. If the submitted command string contains multiple SQL commands, all but the last PGresult are discarded by PQexec.

• PQexec always collects the command’s entire result, buffering it in a single PGresult. While this simplifies error-handling logic for the application, it can be impractical for results containing many rows.

Applications that do not like these limitations can instead use the underlying functions that PQexec is built from: PQsendQuery and PQgetResult. There are also PQsendQueryParams, PQsendPrepare, PQsendQueryPrepared, PQsendDescribePrepared, and PQsendDescribePortal, which can be used with PQgetResult to duplicate the functionality of PQexecParams, PQprepare, PQexecPrepared, PQdescribePrepared, and PQdescribePortal respectively.

PQsendQuery

Submits a command to the server without waiting for the result(s). 1 is returned if the command was successfully dispatched and 0 if not (in which case, use PQerrorMessage to get more information about the failure).
int PQsendQuery(PGconn *conn, const char *command);

After successfully calling PQsendQuery, call PQgetResult one or more times to obtain the results. PQsendQuery cannot be called again (on the same connection) until PQgetResult has returned a null pointer, indicating that the command is done.

PQsendQueryParams

Submits a command and separate parameters to the server without waiting for the result(s).

int PQsendQueryParams(PGconn *conn,
    const char *command,
    int nParams,
    const Oid *paramTypes,
    const char * const *paramValues,
    const int *paramLengths,
    const int *paramFormats,
    int resultFormat);

This is equivalent to PQsendQuery except that query parameters can be specified separately from the query string. The function's parameters are handled identically to PQexecParams. Like PQexecParams, it will not work on 2.0-protocol connections, and it allows only one command in the query string.

PQsendPrepare

Sends a request to create a prepared statement with the given parameters, without waiting for completion.

int PQsendPrepare(PGconn *conn,
    const char *stmtName,
    const char *query,
    int nParams,
    const Oid *paramTypes);

This is an asynchronous version of PQprepare: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to determine whether the server successfully created the prepared statement. The function's parameters are handled identically to PQprepare. Like PQprepare, it will not work on 2.0-protocol connections.

PQsendQueryPrepared

Sends a request to execute a prepared statement with given parameters, without waiting for the result(s).

int PQsendQueryPrepared(PGconn *conn,
    const char *stmtName,
    int nParams,
    const char * const *paramValues,
    const int *paramLengths,
    const int *paramFormats,
    int resultFormat);

This is similar to PQsendQueryParams, but the command to be executed is specified by naming a previously-prepared statement, instead of giving a query string. The function's parameters
are handled identically to PQexecPrepared. Like PQexecPrepared, it will not work on 2.0-protocol connections.

PQsendDescribePrepared
Submits a request to obtain information about the specified prepared statement, without waiting for completion.

int PQsendDescribePrepared(PGconn *conn, const char *stmtName);

This is an asynchronous version of PQdescribePrepared: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to obtain the results. The function's parameters are handled identically to PQdescribePrepared. Like PQdescribePrepared, it will not work on 2.0-protocol connections.

PQsendDescribePortal
Submits a request to obtain information about the specified portal, without waiting for completion.

int PQsendDescribePortal(PGconn *conn, const char *portalName);

This is an asynchronous version of PQdescribePortal: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to obtain the results. The function's parameters are handled identically to PQdescribePortal. Like PQdescribePortal, it will not work on 2.0-protocol connections.

PQgetResult
Waits for the next result from a prior PQsendQuery, PQsendQueryParams, PQsendPrepare, PQsendQueryPrepared, PQsendDescribePrepared, or PQsendDescribePortal call, and returns it. A null pointer is returned when the command is complete and there will be no more results.

PGresult *PQgetResult(PGconn *conn);

PQgetResult must be called repeatedly until it returns a null pointer, indicating that the command is done. (If called when no command is active, PQgetResult will just return a null pointer at once.) Each non-null result from PQgetResult should be processed using the same PGresult accessor functions previously described. Don't forget to free each result object with PQclear when done with it. Note that PQgetResult will block only if a command is active and the necessary response data has not yet been read by PQconsumeInput.

Note
Even when PQresultStatus indicates a fatal error, PQgetResult should be called until it returns a null pointer, to allow libpq to process the error information completely.

Using PQsendQuery and PQgetResult solves one of PQexec's problems: If a command string contains multiple SQL commands, the results of those commands can be obtained individually. (This allows a simple form of overlapped processing, by the way: the client can be handling the results of one command while the server is still working on later queries in the same command string.)

Another frequently-desired feature that can be obtained with PQsendQuery and PQgetResult is retrieving large query results a row at a time. This is discussed in Section 34.5.
By itself, calling PQgetResult will still cause the client to block until the server completes the next SQL command. This can be avoided by proper use of two more functions:

PQconsumeInput

If input is available from the server, consume it.

```c
int PQconsumeInput(PGconn *conn);
```

PQconsumeInput normally returns 1 indicating "no error", but returns 0 if there was some kind of trouble (in which case PQerrorMessage can be consulted). Note that the result does not say whether any input data was actually collected. After calling PQconsumeInput, the application can check PQisBusy and/or PQnotifies to see if their state has changed.

PQconsumeInput can be called even if the application is not prepared to deal with a result or notification just yet. The function will read available data and save it in a buffer, thereby causing a select() read-ready indication to go away. The application can thus use PQconsumeInput to clear the select() condition immediately, and then examine the results at leisure.

PQisBusy

Returns 1 if a command is busy, that is, PQgetResult would block waiting for input. A 0 return indicates that PQgetResult can be called with assurance of not blocking.

```c
int PQisBusy(PGconn *conn);
```

PQisBusy will not itself attempt to read data from the server; therefore PQconsumeInput must be invoked first, or the busy state will never end.

A typical application using these functions will have a main loop that uses select() or poll() to wait for all the conditions that it must respond to. One of the conditions will be input available from the server, which in terms of select() means readable data on the file descriptor identified by PQsocket. When the main loop detects input ready, it should call PQconsumeInput to read the input. It can then call PQisBusy, followed by PQgetResult if PQisBusy returns false (0). It can also call PQnotifies to detect NOTIFY messages (see Section 34.8).

A client that uses PQsendQuery/PQgetResult can also attempt to cancel a command that is still being processed by the server; see Section 34.6. But regardless of the return value of PQcancel, the application must continue with the normal result-reading sequence using PQgetResult. A successful cancellation will simply cause the command to terminate sooner than it would have otherwise.

By using the functions described above, it is possible to avoid blocking while waiting for input from the database server. However, it is still possible that the application will block waiting to send output to the server. This is relatively uncommon but can happen if very long SQL commands or data values are sent. (It is much more probable if the application sends data via COPY IN, however.) To prevent this possibility and achieve completely nonblocking database operation, the following additional functions can be used.

PQsetnonblocking

Sets the nonblocking status of the connection.

```c
int PQsetnonblocking(PGconn *conn, int arg);
```

Sets the state of the connection to nonblocking if arg is 1, or blocking if arg is 0. Returns 0 if OK, -1 if error.
In the nonblocking state, calls to PQsendQuery, PQputline, PQput nbytes, PQput- 
CopyData, and PQendcopy will not block but instead return an error if they need to be called 
again.

Note that PQexec does not honor nonblocking mode; if it is called, it will act in blocking fashion 
anyway.

PQisnonblocking

Returns the blocking status of the database connection.

    int PQisnonblocking(const PGconn *conn);

Returns 1 if the connection is set to nonblocking mode and 0 if blocking.

PQflush

Attempts to flush any queued output data to the server. Returns 0 if successful (or if the send 
queue is empty), -1 if it failed for some reason, or 1 if it was unable to send all the data in the send 
queue yet (this case can only occur if the connection is nonblocking).

    int PQflush(PGconn *conn);

After sending any command or data on a nonblocking connection, call PQflush. If it returns 1, wait 
for the socket to become read- or write-ready. If it becomes write-ready, call PQflush again. If it 
becomes read-ready, call PQconsumeInput, then call PQflush again. Repeat until PQflush re-
turns 0. (It is necessary to check for read-ready and drain the input with PQconsumeInput, because 
the server can block trying to send us data, e.g. NOTICE messages, and won't read our data until we 
read its.) Once PQflush returns 0, wait for the socket to be read-ready and then read the response 
as described above.

### 34.5. Retrieving Query Results Row-By-Row

Ordinarily, libpq collects a SQL command's entire result and returns it to the application as a single 
PGresult. This can be unworkable for commands that return a large number of rows. For such 
cases, applications can use PQsendQuery and PQgetResult in single-row mode. In this mode, 
the result row(s) are returned to the application one at a time, as they are received from the server.

To enter single-row mode, call PQsetSingleRowMode immediately after a successful call of 
PQsendQuery (or a sibling function). This mode selection is effective only for the currently execu-
ting query. Then call PQgetResult repeatedly, until it returns null, as documented in Section 34.4. If 
the query returns any rows, they are returned as individual PGresult objects, which look like nor-
mal query results except for having status code PGRES_SINGLE_TUPLE instead of PGRES_TU-
PLES_OK. After the last row, or immediately if the query returns zero rows, a zero-row object with 
status PGRES_TUPLES_OK is returned; this is the signal that no more rows will arrive. (But note 
that it is still necessary to continue calling PQgetResult until it returns null.) All of these PGre-
sult objects will contain the same row description data (column names, types, etc) that an ordinary 
PGresult object for the query would have. Each object should be freed with PQclear as usual.

PQsetSingleRowMode

Select single-row mode for the currently-executing query.

    int PQsetSingleRowMode(PGconn *conn);

This function can only be called immediately after PQsendQuery or one of its sibling functions, 
before any other operation on the connection such as PQconsumeInput or PQgetResult. If
called at the correct time, the function activates single-row mode for the current query and returns 1. Otherwise the mode stays unchanged and the function returns 0. In any case, the mode reverts to normal after completion of the current query.

Caution

While processing a query, the server may return some rows and then encounter an error, causing the query to be aborted. Ordinarily, libpq discards any such rows and reports only the error. But in single-row mode, those rows will have already been returned to the application. Hence, the application will see some PGRES_SINGLE_TUPLE PGresult objects followed by a PGRES_FATAL_ERROR object. For proper transactional behavior, the application must be designed to discard or undo whatever has been done with the previously-processed rows, if the query ultimately fails.

34.6. Canceling Queries in Progress

A client application can request cancellation of a command that is still being processed by the server, using the functions described in this section.

PQgetCancel

Creates a data structure containing the information needed to cancel a command issued through a particular database connection.

PGcancel *PQgetCancel(PGconn *conn);

PQgetCancel creates a PGcancel object given a PGconn connection object. It will return NULL if the given conn is NULL or an invalid connection. The PGcancel object is an opaque structure that is not meant to be accessed directly by the application; it can only be passed to PQcancel or PQfreeCancel.

PQfreeCancel

Frees a data structure created by PQgetCancel.

void PQfreeCancel(PGcancel *cancel);

PQfreeCancel frees a data object previously created by PQgetCancel.

PQcancel

Requests that the server abandon processing of the current command.

int PQcancel(PGcancel *cancel, char *errbuf, int errbufsize);

The return value is 1 if the cancel request was successfully dispatched and 0 if not. If not, errbuf is filled with an explanatory error message. errbuf must be a char array of size errbufsize (the recommended size is 256 bytes).

Successful dispatch is no guarantee that the request will have any effect, however. If the cancellation is effective, the current command will terminate early and return an error result. If the cancellation fails (say, because the server was already done processing the command), then there will be no visible result at all.
PQcancel can safely be invoked from a signal handler, if the errbuf is a local variable in the signal handler. The PGcancel object is read-only as far as PQcancel is concerned, so it can also be invoked from a thread that is separate from the one manipulating the PGconn object.

PQrequestCancel

PQrequestCancel is a deprecated variant of PQcancel.

int PQrequestCancel(PGconn *conn);

Requests that the server abandon processing of the current command. It operates directly on the PGconn object, and in case of failure stores the error message in the PGconn object (whence it can be retrieved by PQerrorMessage). Although the functionality is the same, this approach creates hazards for multiple-thread programs and signal handlers, since it is possible that overwriting the PGconn's error message will mess up the operation currently in progress on the connection.

34.7. The Fast-Path Interface

PostgreSQL provides a fast-path interface to send simple function calls to the server.

Tip

This interface is somewhat obsolete, as one can achieve similar performance and greater functionality by setting up a prepared statement to define the function call. Then, executing the statement with binary transmission of parameters and results substitutes for a fast-path function call.

The function PQfn requests execution of a server function via the fast-path interface:

PGresult *PQfn(PGconn *conn,
                  int fnid,
                  int *result_buf,
                  int *result_len,
                  int result_is_int,
                  const PQArgBlock *args,
                  int nargs);

typedef struct
{   int len;
    int isint;
    union
    {
        int *ptr;
        int integer;
    } u;
} PQArgBlock;

The fnid argument is the OID of the function to be executed. args and nargs define the parameters to be passed to the function; they must match the declared function argument list. When the isint field of a parameter structure is true, the u.integer value is sent to the server as an integer of the indicated length (this must be 2 or 4 bytes); proper byte-swapping occurs. When isint is false, the indicated number of bytes at *u.ptr are sent with no processing; the data must be in the format expected by the server for binary transmission of the function's argument data type. (The
declaration of u.ptr as being of type int * is historical; it would be better to consider it void *

) result_buf points to the buffer in which to place the function's return value. The caller must have allocated sufficient space to store the return value. (There is no check!) The actual result length in bytes will be returned in the integer pointed to by result_len. If a 2- or 4-byte integer result is expected, set result_is_int to 1, otherwise set it to 0. Setting result_is_int to 1 causes libpq to byte-swap the value if necessary, so that it is delivered as a proper int value for the client machine; note that a 4-byte integer is delivered into *result_buf for either allowed result size. When result_is_int is 0, the binary-format byte string sent by the server is returned unmodified. (In this case it's better to consider result_buf as being of type void *)

PQfn always returns a valid PGresult pointer. The result status should be checked before the result is used. The caller is responsible for freeing the PGresult with PQclear when it is no longer needed.

Note that it is not possible to handle null arguments, null results, nor set-valued results when using this interface.

### 34.8. Asynchronous Notification

PostgreSQL offers asynchronous notification via the LISTEN and NOTIFY commands. A client session registers its interest in a particular notification channel with the LISTEN command (and can stop listening with the UNLISTEN command). All sessions listening on a particular channel will be notified asynchronously when a NOTIFY command with that channel name is executed by any session. A “payload” string can be passed to communicate additional data to the listeners.

libpq applications submit LISTEN, UNLISTEN, and NOTIFY commands as ordinary SQL commands. The arrival of NOTIFY messages can subsequently be detected by calling PQnotifies.

The function PQnotifies returns the next notification from a list of unhandled notification messages received from the server. It returns a null pointer if there are no pending notifications. Once a notification is returned from PQnotifies, it is considered handled and will be removed from the list of notifications.

```c
PGnotify *PQnotifies(PGconn *conn);
```

typedef struct pgNotify
{
    char *relname;              /* notification channel name */
    int  be_pid;                /* process ID of notifying server process */
    char *extra;                /* notification payload string */
} PGnotify;

After processing a PGnotify object returned by PQnotifies, be sure to free it with PQfreemem. It is sufficient to free the PGnotify pointer; the relname and extra fields do not represent separate allocations. (The names of these fields are historical; in particular, channel names need not have anything to do with relation names.)

Example 34.2 gives a sample program that illustrates the use of asynchronous notification.

PQnotifies does not actually read data from the server; it just returns messages previously absorbed by another libpq function. In ancient releases of libpq, the only way to ensure timely receipt of NOTIFY messages was to constantly submit commands, even empty ones, and then check PQnotifies after each PQexec. While this still works, it is deprecated as a waste of processing power.

A better way to check for NOTIFY messages when you have no useful commands to execute is to call PQconsumeInput, then check PQnotifies. You can use select() to wait for data to arrive from the server, thereby using no CPU power unless there is something to do. (See PQsocket
to obtain the file descriptor number to use with select(). Note that this will work OK whether you submit commands with PQsendQuery/PQgetResult or simply use PQexec. You should, however, remember to check PQnotifies after each PQgetResult or PQexec, to see if any notifications came in during the processing of the command.

34.9. Functions Associated with the COPY Command

The COPY command in PostgreSQL has options to read from or write to the network connection used by libpq. The functions described in this section allow applications to take advantage of this capability by supplying or consuming copied data.

The overall process is that the application first issues the SQL COPY command via PQexec or one of the equivalent functions. The response to this (if there is no error in the command) will be a PGresult object bearing a status code of PGRES_COPY_OUT or PGRES_COPY_IN (depending on the specified copy direction). The application should then use the functions of this section to receive or transmit data rows. When the data transfer is complete, another PGresult object is returned to indicate success or failure of the transfer. Its status will be PGRES_COMMAND_OK for success or PGRES_FATAL_ERROR if some problem was encountered. At this point further SQL commands can be issued via PQexec. (It is not possible to execute other SQL commands using the same connection while the COPY operation is in progress.)

If a COPY command is issued via PQexec in a string that could contain additional commands, the application must continue fetching results via PQgetResult after completing the COPY sequence. Only when PQgetResult returns NULL is it certain that the PQexec command string is done and it is safe to issue more commands.

The functions of this section should be executed only after obtaining a result status of PGRES_COPY_OUT or PGRES_COPY_IN from PQexec or PQgetResult.

A PGresult object bearing one of these status values carries some additional data about the COPY operation that is starting. This additional data is available using functions that are also used in connection with query results:

PQnfields

Returns the number of columns (fields) to be copied.

PQbinaryTuples

0 indicates the overall copy format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary. See COPY for more information.

PQfformat

Returns the format code (0 for text, 1 for binary) associated with each column of the copy operation. The per-column format codes will always be zero when the overall copy format is textual, but the binary format can support both text and binary columns. (However, as of the current implementation of COPY, only binary columns appear in a binary copy; so the per-column formats always match the overall format at present.)

Note

These additional data values are only available when using protocol 3.0. When using protocol 2.0, all these functions will return 0.
34.9.1. Functions for Sending COPY Data

These functions are used to send data during COPY FROM STDIN. They will fail if called when the connection is not in COPY_IN state.

PQputCopyData

Sends data to the server during COPY_IN state.

```c
int PQputCopyData(PGconn *conn, const char *buffer, int nbytes);
```

Transmits the COPY data in the specified `buffer`, of length `nbytes`, to the server. The result is 1 if the data was queued, zero if it was not queued because of full buffers (this will only happen in nonblocking mode), or -1 if an error occurred. (Use PQerrorMessage to retrieve details if the return value is -1. If the value is zero, wait for write-ready and try again.)

The application can divide the COPY data stream into buffer loads of any convenient size. Buffer-load boundaries have no semantic significance when sending. The contents of the data stream must match the data format expected by the COPY command; see COPY for details.

PQputCopyEnd

Sends end-of-data indication to the server during COPY_IN state.

```c
int PQputCopyEnd(PGconn *conn, const char *errormsg);
```

Ends the COPY_IN operation successfully if `errormsg` is NULL. If `errormsg` is not NULL then the COPY is forced to fail, with the string pointed to by `errormsg` used as the error message. (One should not assume that this exact error message will come back from the server, however, as the server might have already failed the COPY for its own reasons. Also note that the option to force failure does not work when using pre-3.0-protocol connections.)

The result is 1 if the termination message was sent; or in nonblocking mode, this may only indicate that the termination message was successfully queued. (In nonblocking mode, to be certain that the data has been sent, you should next wait for write-ready and call PQflush, repeating until it returns zero.) Zero indicates that the function could not queue the termination message because of full buffers; this will only happen in nonblocking mode. (In this case, wait for write-ready and try the PQputCopyEnd call again.) If a hard error occurs, -1 is returned; you can use PQerrorMessage to retrieve details.

After successfully calling PQputCopyEnd, call PQgetResult to obtain the final result status of the COPY command. One can wait for this result to be available in the usual way. Then return to normal operation.

34.9.2. Functions for Receiving COPY Data

These functions are used to receive data during COPY TO STDOUT. They will fail if called when the connection is not in COPY_OUT state.

PQgetCopyData

Receives data from the server during COPY_OUT state.
int PQgetCopyData(PGconn *conn,  
    char **buffer,  
    int async);

Attempts to obtain another row of data from the server during a `COPY`. Data is always returned one data row at a time; if only a partial row is available, it is not returned. Successful return of a data row involves allocating a chunk of memory to hold the data. The `buffer` parameter must be non-NULL. `*buffer` is set to point to the allocated memory, or to NULL in cases where no buffer is returned. A non-NULL result buffer should be freed using `PQfreemem` when no longer needed.

When a row is successfully returned, the return value is the number of data bytes in the row (this will always be greater than zero). The returned string is always null-terminated, though this is probably only useful for textual `COPY`. A result of zero indicates that the `COPY` is still in progress, but no row is yet available (this is only possible when async is true). A result of -1 indicates that the `COPY` is done. A result of -2 indicates that an error occurred (consult `PQerrorMessage` for the reason).

When `async` is true (not zero), `PQgetCopyData` will not block waiting for input; it will return zero if the `COPY` is still in progress but no complete row is available. (In this case wait for read-ready and then call `PQconsumeInput` before calling `PQgetCopyData` again.) When `async` is false (zero), `PQgetCopyData` will block until data is available or the operation completes.

After `PQgetCopyData` returns -1, call `PQgetResult` to obtain the final result status of the `COPY` command. One can wait for this result to be available in the usual way. Then return to normal operation.

### 34.9.3. Obsolete Functions for `COPY`

These functions represent older methods of handling `COPY`. Although they still work, they are deprecated due to poor error handling, inconvenient methods of detecting end-of-data, and lack of support for binary or nonblocking transfers.

**PQgetline**

Reads a newline-terminated line of characters (transmitted by the server) into a buffer string of size `length`.

```
int PQgetline(PGconn *conn,  
    char *buffer,  
    int length);
```

This function copies up to `length`-1 characters into the buffer and converts the terminating newline into a zero byte. `PQgetline` returns EOF at the end of input, 0 if the entire line has been read, and 1 if the buffer is full but the terminating newline has not yet been read.

Note that the application must check to see if a new line consists of the two characters \\
which indicates that the server has finished sending the results of the `COPY` command. If the application might receive lines that are more than `length`-1 characters long, care is needed to be sure it recognizes the \
line correctly (and does not, for example, mistake the end of a long data line for a terminator line).

**PQgetlineAsync**

Reads a row of `COPY` data (transmitted by the server) into a buffer without blocking.

```
int PQgetlineAsync(PGconn *conn,  
    char *buffer,
```
This function is similar to PQgetline, but it can be used by applications that must read COPY data asynchronously, that is, without blocking. Having issued the COPY command and gotten a PGRES_COPY_OUT response, the application should call PQconsumeInput and PQgetlineAsync until the end-of-data signal is detected.

Unlike PQgetline, this function takes responsibility for detecting end-of-data.

On each call, PQgetlineAsync will return data if a complete data row is available in libpq's input buffer. Otherwise, no data is returned until the rest of the row arrives. The function returns -1 if the end-of-copy-data marker has been recognized, or 0 if no data is available, or a positive number giving the number of bytes of data returned. If -1 is returned, the caller must next call PQendcopy, and then return to normal processing.

The data returned will not extend beyond a data-row boundary. If possible a whole row will be returned at one time. But if the buffer offered by the caller is too small to hold a row sent by the server, then a partial data row will be returned. With textual data this can be detected by testing whether the last returned byte is \n or not. (In a binary COPY, actual parsing of the COPY data format will be needed to make the equivalent determination.) The returned string is not null-terminated. (If you want to add a terminating null, be sure to pass a bufsize one smaller than the room actually available.)

PQputline

Sends a null-terminated string to the server. Returns 0 if OK and EOF if unable to send the string.

int PQputline(PGconn *conn,
               const char *string);

The COPY data stream sent by a series of calls to PQputline has the same format as that returned by PQgetlineAsync, except that applications are not obliged to send exactly one data row per PQputline call; it is okay to send a partial line or multiple lines per call.

Note

Before PostgreSQL protocol 3.0, it was necessary for the application to explicitly send the two characters \. as a final line to indicate to the server that it had finished sending COPY data. While this still works, it is deprecated and the special meaning of \. can be expected to be removed in a future release. It is sufficient to call PQendcopy after having sent the actual data.

PQputnbytes

Sends a non-null-terminated string to the server. Returns 0 if OK and EOF if unable to send the string.

int PQputnbytes(PGconn *conn,
                const char *buffer,
                int nbytes);

This is exactly like PQputline, except that the data buffer need not be null-terminated since the number of bytes to send is specified directly. Use this procedure when sending binary data.

PQendcopy

Synchronizes with the server.
int PQendcopy(PGconn *conn);

This function waits until the server has finished the copying. It should either be issued when the last string has been sent to the server using PQputline or when the last string has been received from the server using PGgetline. It must be issued or the server will get “out of sync” with the client. Upon return from this function, the server is ready to receive the next SQL command. The return value is 0 on successful completion, nonzero otherwise. (Use PQerrorMessage to retrieve details if the return value is nonzero.)

When using PQgetResult, the application should respond to a PGRES_COPY_OUT result by executing PQgetline repeatedly, followed by PQendcopy after the terminator line is seen. It should then return to the PQgetResult loop until PQgetResult returns a null pointer. Similarly a PGRES_COPY_IN result is processed by a series of PQputline calls followed by PQendcopy, then return to the PQgetResult loop. This arrangement will ensure that a COPY command embedded in a series of SQL commands will be executed correctly.

Older applications are likely to submit a COPY via PQexec and assume that the transaction is done after PQendcopy. This will work correctly only if the COPY is the only SQL command in the command string.

### 34.10. Control Functions

These functions control miscellaneous details of libpq’s behavior.

**PQclientEncoding**

Returns the client encoding.

```c
int PQclientEncoding(const PGconn *conn);
```

Note that it returns the encoding ID, not a symbolic string such as EUC_JP. If unsuccessful, it returns -1. To convert an encoding ID to an encoding name, you can use:

```c
char *pg_encoding_to_char(int encoding_id);
```

**PQsetClientEncoding**

Sets the client encoding.

```c
int PQsetClientEncoding(PGconn *conn, const char *encoding);
```

`conn` is a connection to the server, and `encoding` is the encoding you want to use. If the function successfully sets the encoding, it returns 0, otherwise -1. The current encoding for this connection can be determined by using PQclientEncoding.

**PQsetErrorVerbosity**

Determines the verbosity of messages returned by PQerrorMessage and PQresultErrorMessage.

```c
typedef enum {
  PQERRORS_TERSE,
  PQERRORS_DEFAULT,
} PQERRORS;
```
PQERRORS_VERBOSE
} PGVerbosity;

PGVerbosity PQsetErrorVerbosity(PGconn *conn, PGVerbosity verbosity);

PQsetErrorVerbosity sets the verbosity mode, returning the connection’s previous setting. In **TERSE** mode, returned messages include severity, primary text, and position only; this will normally fit on a single line. The default mode produces messages that include the above plus any detail, hint, or context fields (these might span multiple lines). The **VERBOSE** mode includes all available fields. Changing the verbosity does not affect the messages available from already-existing PGresult objects, only subsequently-created ones. (But see PQresultVerboseErrorMessage if you want to print a previous error with a different verbosity.)

PQsetErrorContextVisibility

Determines the handling of CONTEXT fields in messages returned by PQerrorMessage and PQresultErrorMessage.

typedef enum
{
    PQSHOW_CONTEXT_NEVER,
    PQSHOW_CONTEXT_ERRORS,
    PQSHOW_CONTEXT_ALWAYS
} PGContextVisibility;

PGContextVisibility PQsetErrorContextVisibility(PGconn *conn, PGContextVisibility show_context);

PQsetErrorContextVisibility sets the context display mode, returning the connection’s previous setting. This mode controls whether the CONTEXT field is included in messages (unless the verbosity setting is **TERSE**, in which case CONTEXT is never shown). The **NEVER** mode never includes CONTEXT, while **ALWAYS** always includes it if available. In **ERRORS** mode (the default), CONTEXT fields are included only for error messages, not for notices and warnings. Changing this mode does not affect the messages available from already-existing PGresult objects, only subsequently-created ones. (But see PQresultVerboseErrorMessage if you want to print a previous error with a different display mode.)

PQtrace

Enables tracing of the client/server communication to a debugging file stream.

void PQtrace(PGconn *conn, FILE *stream);

Note

On Windows, if the libpq library and an application are compiled with different flags, this function call will crash the application because the internal representation of the FILE pointers differ. Specifically, multithreaded/single-threaded, release/debug, and static/dynamic flags should be the same for the library and all applications using that library.

PQuntrace

Disables tracing started by PQtrace.
void PQuntrace(PGconn *conn);

34.11. Miscellaneous Functions

As always, there are some functions that just don’t fit anywhere.

PQfreemem
Frees memory allocated by libpq.

void PQfreemem(void *ptr);
Frees memory allocated by libpq, particularly PQescapeByteaConn, PQescapeBytea, PQunescapeBytea, and PQnotifies. It is particularly important that this function, rather than free(), be used on Microsoft Windows. This is because allocating memory in a DLL and releasing it in the application works only if multithreaded/single-threaded, release/debug, and static/dynamic flags are the same for the DLL and the application. On non-Microsoft Windows platforms, this function is the same as the standard library function free().

PQconninfoFree
Frees the data structures allocated by PQconndefaults or PQconninfoParse.

void PQconninfoFree(PQconninfoOption *connOptions);
A simple PQfreemem will not do for this, since the array contains references to subsidiary strings.

PQencryptPasswordConn
Prepares the encrypted form of a PostgreSQL password.

char *PQencryptPasswordConn(PGconn *conn, const char *passwd, const char *user, const char *algorithm);

This function is intended to be used by client applications that wish to send commands like ALTER USER joe PASSWORD 'pwd'. It is good practice not to send the original cleartext password in such a command, because it might be exposed in command logs, activity displays, and so on. Instead, use this function to convert the password to encrypted form before it is sent.

The passwd and user arguments are the cleartext password, and the SQL name of the user it is for. algorithm specifies the encryption algorithm to use to encrypt the password. Currently supported algorithms are md5 and scram-sha-256 (on and off are also accepted as aliases for md5, for compatibility with older server versions). Note that support for scram-sha-256 was introduced in PostgreSQL version 10, and will not work correctly with older server versions. If algorithm is NULL, this function will query the server for the current value of the password_encryption setting. That can block, and will fail if the current transaction is aborted, or if the connection is busy executing another query. If you wish to use the default algorithm for the server but want to avoid blocking, query password_encryption yourself before calling PQencryptPasswordConn, and pass that value as the algorithm.

The return value is a string allocated by malloc. The caller can assume the string doesn’t contain any special characters that would require escaping. Use PQfreemem to free the result when done with it. On error, returns NULL, and a suitable message is stored in the connection object.
PQencryptPassword
Prepares the md5-encrypted form of a PostgreSQL password.

char *PQencryptPassword(const char *passwd, const char *user);
PQencryptPassword is an older, deprecated version of PQencryptPasswordConn. The difference is that PQencryptPassword does not require a connection object, and md5 is always used as the encryption algorithm.

PQmakeEmptyPGresult
Constructs an empty PGresult object with the given status.

PGresult *PQmakeEmptyPGresult(PGconn *conn, ExecStatusType status);
This is libpq's internal function to allocate and initialize an empty PGresult object. This function returns NULL if memory could not be allocated. It is exported because some applications find it useful to generate result objects (particularly objects with error status) themselves. If conn is not null and status indicates an error, the current error message of the specified connection is copied into the PGresult. Also, if conn is not null, any event procedures registered in the connection are copied into the PGresult. (They do not get PGEVT_RESULTCREATE calls, but see PQfireResultCreateEvents.) Note that PQclear should eventually be called on the object, just as with a PGresult returned by libpq itself.

PQfireResultCreateEvents
Fires a PGEVT_RESULTCREATE event (see Section 34.13) for each event procedure registered in the PGresult object. Returns non-zero for success, zero if any event procedure fails.

int PQfireResultCreateEvents(PGconn *conn, PGresult *res);
The conn argument is passed through to event procedures but not used directly. It can be NULL if the event procedures won't use it.
Event procedures that have already received a PGEVT_RESULTCREATE or PGEVT_RESULTCOPY event for this object are not fired again.

The main reason that this function is separate from PQmakeEmptyPGresult is that it is often appropriate to create a PGresult and fill it with data before invoking the event procedures.

PQcopyResult
Makes a copy of a PGresult object. The copy is not linked to the source result in any way and PQclear must be called when the copy is no longer needed. If the function fails, NULL is returned.

PGresult *PQcopyResult(const PGresult *src, int flags);
This is not intended to make an exact copy. The returned result is always put into PGRES_TUPLES_OK status, and does not copy any error message in the source. (It does copy the command status string, however.) The flags argument determines what else is copied. It is a bitwise OR of several flags. PG_COPYRES_ATTRS specifies copying the source result's attributes (column definitions). PG_COPYRES_TUPLES specifies copying the source result's tuples. (This implies copying the attributes, too.) PG_COPYRES_NOTICEHOOKS specifies copying the source result's
notify hooks. `PG_COPYRES_EVENTS` specifies copying the source result's events. (But any instance data associated with the source is not copied.)

**PQsetResultAttrs**

Sets the attributes of a `PGresult` object.

```c
int PQsetResultAttrs(PGresult *res, int numAttributes, PGresAttDesc *attDescs);
```

The provided `attDescs` are copied into the result. If the `attDescs` pointer is NULL or `numAttributes` is less than one, the request is ignored and the function succeeds. If `res` already contains attributes, the function will fail. If the function fails, the return value is zero. If the function succeeds, the return value is non-zero.

**PQsetvalue**

Sets a tuple field value of a `PGresult` object.

```c
int PQsetvalue(PGresult *res, int tup_num, int field_num, char *value, int len);
```

The function will automatically grow the result's internal tuples array as needed. However, the `tup_num` argument must be less than or equal to `PQntuples`, meaning this function can only grow the tuples array one tuple at a time. But any field of any existing tuple can be modified in any order. If a value at `field_num` already exists, it will be overwritten. If `len` is -1 or `value` is NULL, the field value will be set to an SQL null value. The value is copied into the result's private storage, thus is no longer needed after the function returns. If the function fails, the return value is zero. If the function succeeds, the return value is non-zero.

**PQresultAlloc**

Allocate subsidiary storage for a `PGresult` object.

```c
void *PQresultAlloc(PGresult *res, size_t nBytes);
```

Any memory allocated with this function will be freed when `res` is cleared. If the function fails, the return value is NULL. The result is guaranteed to be adequately aligned for any type of data, just as for `malloc`.

**PQlibVersion**

Return the version of libpq that is being used.

```c
int PQlibVersion(void);
```

The result of this function can be used to determine, at run time, whether specific functionality is available in the currently loaded version of libpq. The function can be used, for example, to determine which connection options are available in `PQconnectdb`.

The result is formed by multiplying the library's major version number by 10000 and adding the minor version number. For example, version 10.1 will be returned as 100001, and version 11.0 will be returned as 110000.

Prior to major version 10, PostgreSQL used three-part version numbers in which the first two parts together represented the major version. For those versions, `PQlibVersion` uses two digits for each part; for example version 9.1.5 will be returned as 90105, and version 9.2.0 will be returned as 90200.
Therefore, for purposes of determining feature compatibility, applications should divide the result of `PQlibVersion()` by 100, not 10000, to determine a logical major version number. In all release series, only the last two digits differ between minor releases (bug-fix releases).

**Note**

This function appeared in PostgreSQL version 9.1, so it cannot be used to detect required functionality in earlier versions, since calling it will create a link dependency on version 9.1 or later.

### 34.12. Notice Processing

Notice and warning messages generated by the server are not returned by the query execution functions, since they do not imply failure of the query. Instead they are passed to a notice handling function, and execution continues normally after the handler returns. The default notice handling function prints the message on `stderr`, but the application can override this behavior by supplying its own handling function.

For historical reasons, there are two levels of notice handling, called the notice receiver and notice processor. The default behavior is for the notice receiver to format the notice and pass a string to the notice processor for printing. However, an application that chooses to provide its own notice receiver will typically ignore the notice processor layer and just do all the work in the notice receiver.

The function `PQsetNoticeReceiver()` sets or examines the current notice receiver for a connection object. Similarly, `PQsetNoticeProcessor()` sets or examines the current notice processor.

```c
typedef void (*PQnoticeReceiver) (void *arg, const PGresult *res);
PQnoticeReceiver
PQsetNoticeReceiver(PGconn *conn,
                    PQnoticeReceiver proc,
                    void *arg);

typedef void (*PQnoticeProcessor) (void *arg, const char *message);
PQnoticeProcessor
PQsetNoticeProcessor(PGconn *conn,
                      PQnoticeProcessor proc,
                      void *arg);
```

Each of these functions returns the previous notice receiver or processor function pointer, and sets the new value. If you supply a null function pointer, no action is taken, but the current pointer is returned.

When a notice or warning message is received from the server, or generated internally by libpq, the notice receiver function is called. It is passed the message in the form of a `PGRES_NONFATAL_ERROR` `PGresult`. (This allows the receiver to extract individual fields using `PQresultErrorField`, or obtain a complete preformatted message using `PQresultErrorMessage` or `PQresultVerboseErrorMessage`.) The same void pointer passed to `PQsetNoticeReceiver` is also passed. (This pointer can be used to access application-specific state if needed.)

The default notice receiver simply extracts the message (using `PQresultErrorMessage`) and passes it to the notice processor.

The notice processor is responsible for handling a notice or warning message given in text form. It is passed the string text of the message (including a trailing newline), plus a void pointer that is the same
one passed to PQsetNoticeProcessor. (This pointer can be used to access application-specific state if needed.)

The default notice processor is simply:

```c
static void
defaultNoticeProcessor(void *arg, const char *message)
{
    fprintf(stderr, "%s", message);
}
```

Once you have set a notice receiver or processor, you should expect that that function could be called as long as either the PGconn object or PGresult objects made from it exist. At creation of a PGresult, the PGconn's current notice handling pointers are copied into the PGresult for possible use by functions like PQgetvalue.

### 34.13. Event System

libpq's event system is designed to notify registered event handlers about interesting libpq events, such as the creation or destruction of PGconn and PGresult objects. A principal use case is that this allows applications to associate their own data with a PGconn or PGresult and ensure that that data is freed at an appropriate time.

Each registered event handler is associated with two pieces of data, known to libpq only as opaque void * pointers. There is a passthrough pointer that is provided by the application when the event handler is registered with a PGconn. The passthrough pointer never changes for the life of the PGconn and all PGresults generated from it; so if used, it must point to long-lived data. In addition there is an instance data pointer, which starts out NULL in every PGconn and PGresult. This pointer can be manipulated using the PQinstanceData, PQsetInstanceData, PQresultInstanceData and PQsetResultInstanceData functions. Note that unlike the passthrough pointer, instance data of a PGconn is not automatically inherited by PGresults created from it. libpq does not know what passthrough and instance data pointers point to (if anything) and will never attempt to free them — that is the responsibility of the event handler.

### 34.13.1. Event Types

The enum PGEventId names the types of events handled by the event system. All its values have names beginning with PGEVT. For each event type, there is a corresponding event info structure that carries the parameters passed to the event handlers. The event types are:

**PGEVT_REGISTER**

The register event occurs when PQregisterEventProc is called. It is the ideal time to initialize any instanceData an event procedure may need. Only one register event will be fired per event handler per connection. If the event procedure fails, the registration is aborted.

```c
typedef struct
{
    PGconn *conn;
} PGEVT_Register;
```

When a PGEVT_REGISTER event is received, the evtInfo pointer should be cast to a PGEVT_Register *. This structure contains a PGconn that should be in the CONNECTION_OK status, guaranteed if one calls PQregisterEventProc right after obtaining a good PGconn. When returning a failure code, all cleanup must be performed as no PGEVT_CON-NDESTROY event will be sent.
PGEVT_CONNRESET

The connection reset event is fired on completion of PQreset or PQresetPoll. In both cases, the event is only fired if the reset was successful. If the event procedure fails, the entire connection reset will fail; the PGconn is put into CONNECTION_BAD status and PQresetPoll will return PGRES_POLLING_FAILED.

typedef struct
{
   PGconn *conn;
} PGEventConnReset;

When a PGEVT_CONNRESET event is received, the evtInfo pointer should be cast to a PGEventConnReset * . Although the contained PGconn was just reset, all event data remains unchanged. This event should be used to reset/reload/requery any associated instanceData. Note that even if the event procedure fails to process PGEVT_CONNRESET, it will still receive a PGEVT_CONNDESTROY event when the connection is closed.

PGEVT_CONNDESTROY

The connection destroy event is fired in response to PQfinish. It is the event procedure's responsibility to properly clean up its event data as libpq has no ability to manage this memory. Failure to clean up will lead to memory leaks.

typedef struct
{
   PGconn *conn;
} PGEventConnDestroy;

When a PGEVT_CONNDESTROY event is received, the evtInfo pointer should be cast to a PGEventConnDestroy *. This event is fired prior to PQfinish performing any other cleanup. The return value of the event procedure is ignored since there is no way of indicating a failure from PQfinish. Also, an event procedure failure should not abort the process of cleaning up unwanted memory.

PGEVT_RESULTCREATE

The result creation event is fired in response to any query execution function that generates a result, including PQgetResult. This event will only be fired after the result has been created successfully.

typedef struct
{
   PGconn *conn;
   PGresult *result;
} PGEventResultCreate;

When a PGEVT_RESULTCREATE event is received, the evtInfo pointer should be cast to a PGEventResultCreate *. The conn is the connection used to generate the result. This is the ideal place to initialize any instanceData that needs to be associated with the result. If the event procedure fails, the result will be cleared and the failure will be propagated. The event procedure must not try to PQclear the result object for itself. When returning a failure code, all cleanup must be performed as no PGEVT_RESULTDESTROY event will be sent.

PGEVT_RESULTCOPY

The result copy event is fired in response to PQcopyResult. This event will only be fired after the copy is complete. Only event procedures that have successfully handled the
PGEVT_RESULTCREATE or PGEVT_RESULTCOPY event for the source result will receive PGEVT_RESULTCOPY events.

typedef struct
{
    const PGresult *src;
    PGresult *dest;
} PGEventResultCopy;

When a PGEVT_RESULTCOPY event is received, the evtInfo pointer should be cast to a PGEventResultCopy *. The src result is what was copied while the dest result is the copy destination. This event can be used to provide a deep copy of instanceData, since PQ-copyResult cannot do that. If the event procedure fails, the entire copy operation will fail and the dest result will be cleared. When returning a failure code, all cleanup must be performed as no PGEVT_RESULTDESTROY event will be sent for the destination result.

PGEVT_RESULTDESTROY

The result destroy event is fired in response to a PQclear. It is the event procedure’s responsibility to properly clean up its event data as libpq has no ability to manage this memory. Failure to clean up will lead to memory leaks.

typedef struct
{
    PGresult *result;
} PGEventResultDestroy;

When a PGEVT_RESULTDESTROY event is received, the evtInfo pointer should be cast to a PGEventResultDestroy *. This event is fired prior to PQclear performing any other cleanup. The return value of the event procedure is ignored since there is no way of indicating a failure from PQclear. Also, an event procedure failure should not abort the process of cleaning up unwanted memory.

### 34.13.2. Event Callback Procedure

**PGEventProc**

PGEventProc is a typedef for a pointer to an event procedure, that is, the user callback function that receives events from libpq. The signature of an event procedure must be

```c
int eventproc(PGEventId evtId, void *evtInfo, void *passThrough)
```

The `evtId` parameter indicates which PGEVT event occurred. The `evtInfo` pointer must be cast to the appropriate structure type to obtain further information about the event. The `passThrough` parameter is the pointer provided to PQregisterEventProc when the event procedure was registered. The function should return a non-zero value if it succeeds and zero if it fails.

A particular event procedure can be registered only once in any PGconn. This is because the address of the procedure is used as a lookup key to identify the associated instance data.

**Caution**

On Windows, functions can have two different addresses: one visible from outside a DLL and another visible from inside the DLL. One should be careful that only one of these addresses is used with libpq’s event-procedure functions, else confu-
34.13.3. Event Support Functions

PQregisterEventProc

Registers an event callback procedure with libpq.

```c
int PQregisterEventProc(PGconn *conn, PGEventProc proc,
    const char *name, void *passThrough);
```

An event procedure must be registered once on each PGconn you want to receive events about. There is no limit, other than memory, on the number of event procedures that can be registered with a connection. The function returns a non-zero value if it succeeds and zero if it fails.

The proc argument will be called when a libpq event is fired. Its memory address is also used to lookup instanceData. The name argument is used to refer to the event procedure in error messages. This value cannot be NULL or a zero-length string. The name string is copied into the PGconn, so what is passed need not be long-lived. The passThrough pointer is passed to the proc whenever an event occurs. This argument can be NULL.

PQsetInstanceData

Sets the connection conn's instanceData for procedure proc to data. This returns non-zero for success and zero for failure. (Failure is only possible if proc has not been properly registered in conn.)

```c
int PQsetInstanceData(PGconn *conn, PGEventProc proc, void *data);
```

PQinstanceData

Returns the connection conn's instanceData associated with procedure proc, or NULL if there is none.

```c
void *PQinstanceData(const PGconn *conn, PGEventProc proc);
```

PQresultSetInstanceData

Sets the result's instanceData for proc to data. This returns non-zero for success and zero for failure. (Failure is only possible if proc has not been properly registered in the result.)

```c
int PQresultSetInstanceData(PGresult *res, PGEventProc proc,
    void *data);
```

PQresultInstanceData

Returns the result's instanceData associated with proc, or NULL if there is none.

```c
void *PQresultInstanceData(const PGresult *res, PGEventProc proc);
```
34.13.4. Event Example

Here is a skeleton example of managing private data associated with libpq connections and results.

```c
/* required header for libpq events (note: includes libpq-fe.h) */
#include <libpq-events.h>

/* The instanceData */
typedef struct
{
    int n;
    char *str;
} mydata;

/* PGEventProc */
static int myEventProc(PGEventId evtId, void *evtInfo, void *passThrough);

int main(void)
{
    mydata *data;
    PGresult *res;
    PGconn *conn = PQconnectdb("dbname=postgres options=-csearch_path=");

    if (PQstatus(conn) != CONNECTION_OK)
    {
        fprintf(stderr, "Connection to database failed: %s",
            PQerrorMessage(conn));
        PQfinish(conn);
        return 1;
    }

    /* called once on any connection that should receive events. *
    * Sends a PGEVT_REGISTER to myEventProc. */
    if (!PQregisterEventProc(conn, myEventProc, "mydata_proc", NULL))
    {
        fprintf(stderr, "Cannot register PGEventProc\n");
        PQfinish(conn);
        return 1;
    }

    /* conn instanceData is available */
    data = PQinstanceData(conn, myEventProc);

    /* Sends a PGEVT_RESULTCREATE to myEventProc */
    res = PQexec(conn, "SELECT 1 + 1");

    /* result instanceData is available */
    data = PQresultInstanceData(res, myEventProc);

    /* If PG_COPYRES_EVENTS is used, sends a PGEVT_RESULTCOPY to myEventProc */
```
res_copy = PQcopyResult(res, PG_COPYRES_TUPLES | PG_COPYRES_EVENTS);

/* result instanceData is available if PG_COPYRES_EVENTS was used during the PQcopyResult call. */
data = PQresultInstanceData(res_copy, myEventProc);

/* Both clears send a PGEVT_RESULTDESTROY to myEventProc */
PQclear(res);
PQclear(res_copy);

/* Sends a PGEVT_CONNDESTROY to myEventProc */
PQfinish(conn);

return 0;
}

static int
myEventProc(PGEventId evtId, void *evtInfo, void *passThrough)
{
    switch (evtId)
    {
    case PGEVT_REGISTER:
        {
            PGEventRegister *e = (PGEventRegister *)evtInfo;
            mydata *data = get_mydata(e->conn);

            /* associate app specific data with connection */
PQsetInstanceData(e->conn, myEventProc, data);
            break;
        }
    case PGEVT_CONNRESET:
        {
            PGEventConnReset *e = (PGEventConnReset *)evtInfo;
            mydata *data = PQinstanceData(e->conn, myEventProc);

            if (data)
                memset(data, 0, sizeof(mydata));
            break;
        }
    case PGEVT_CONNDESTROY:
        {
            PGEventConnDestroy *e = (PGEventConnDestroy *)evtInfo;
            mydata *data = PQinstanceData(e->conn, myEventProc);

            /* free instance data because the conn is being destroyed */
            if (data)
                free_mydata(data);
            break;
        }
    case PGEVT_RESULTCREATE:
        {

    828
PGEVENTRESULTCREATE *e = (PGEVENTRESULTCREATE *)evtInfo;
mydata *conn_data = PQinstanceData(e->conn,
myEventProc);
mydata *res_data = dup_mydata(conn_data);

/* associate app specific data with result (copy it
from conn) */
PQsetResultInstanceData(e->result, myEventProc,
res_data);
break;
}

case PGEVT_RESULTCOPY:  
{
    PGEVENTRESULTCOPY *e = (PGEVENTRESULTCOPY *)evtInfo;
    mydata *src_data = PQresultInstanceData(e->src,
myEventProc);
    mydata *dest_data = dup_mydata(src_data);

    /* associate app specific data with result (copy it
    from a result) */
PQsetResultInstanceData(e->dest, myEventProc,
dest_data);
    break;
}

case PGEVT_RESULTDESTROY:
{
    PGEVENTRESULTDESTROY *e = (PGEVENTRESULTDESTROY *)evtInfo;
    mydata *data = PQresultInstanceData(e->result,
myEventProc);

    /* free instance data because the result is being
    destroyed */
    if (data)
        free_mydata(data);
    break;
}

/* unknown event ID, just return true. */
default:
    break;
}

return true; /* event processing succeeded */

34.14. Environment Variables

The following environment variables can be used to select default connection parameter values, which
will be used by PQconnectdb, PQsetdbLogin and PQsetdb if no value is directly specified by
the calling code. These are useful to avoid hard-coding database connection information into simple
client applications, for example.

- PGHOST behaves the same as the host connection parameter.
• `PGHOSTADDR` behaves the same as the `hostaddr` connection parameter. This can be set instead of or in addition to `PGHOST` to avoid DNS lookup overhead.

• `PGPORT` behaves the same as the `port` connection parameter.

• `PGDATABASE` behaves the same as the `dbname` connection parameter.

• `PGUSER` behaves the same as the `user` connection parameter.

• `PGPASSWORD` behaves the same as the `password` connection parameter. Use of this environment variable is not recommended for security reasons, as some operating systems allow non-root users to see process environment variables via `ps`; instead consider using a password file (see Section 34.15).

• `PGPASSFILE` behaves the same as the `passfile` connection parameter.

• `PGSERVICE` behaves the same as the `service` connection parameter.

• `PGSERVICEFILE` specifies the name of the per-user connection service file. If not set, it defaults to `~/.pg_service.conf` (see Section 34.16).

• `PGOPTIONS` behaves the same as the `options` connection parameter.

• `PGAPPNAME` behaves the same as the `application_name` connection parameter.

• `PGREQUIRESSL` behaves the same as the `requiressl` connection parameter. This environment variable is deprecated in favor of the `PGSSLMODE` variable; setting both variables suppresses the effect of this one.

• `PGSSLMODE` behaves the same as the `sslmode` connection parameter.

• `PGREQUIREPEER` behaves the same as the `requirepeer` connection parameter.

• `PGKRBSSRVNAME` behaves the same as the `krbsrvname` connection parameter.

• `PGGSSLIB` behaves the same as the `gsslib` connection parameter.

• `PGCONNECT_TIMEOUT` behaves the same as the `connect_timeout` connection parameter.

• `PGCLIENTENCODING` behaves the same as the `client_encoding` connection parameter.

• `PGTARGETSESSIONATTRS` behaves the same as the `target_session_attrs` connection parameter.

The following environment variables can be used to specify default behavior for each PostgreSQL session. (See also the `ALTER ROLE` and `ALTER DATABASE` commands for ways to set default behavior on a per-user or per-database basis.)

• `PGDATESTYLE` sets the default style of date/time representation. (Equivalent to `SET datestyle TO ....`)  

• `PGTZ` sets the default time zone. (Equivalent to `SET timezone TO ....`)  

• `PGGEQO` sets the default mode for the genetic query optimizer. (Equivalent to `SET geqo TO ....`)  

Refer to the SQL command `SET` for information on correct values for these environment variables.
The following environment variables determine internal behavior of libpq; they override compiled-in defaults.

- **PGSYSConfDIR** sets the directory containing the `pg_service.conf` file and in a future version possibly other system-wide configuration files.
- **PGLocaLEDIR** sets the directory containing the `locale` files for message localization.

## 34.15. The Password File

The file `.pgpass` in a user's home directory can contain passwords to be used if the connection requires a password (and no password has been specified otherwise). On Microsoft Windows the file is named `%APPDATA%\postgresql\pgpass.conf` (where `%APPDATA%` refers to the Application Data subdirectory in the user's profile). Alternatively, a password file can be specified using the connection parameter `passfile` or the environment variable `PGPASSFILE`. This file should contain lines of the following format:

```
hostname:port:database:username:password
```

(You can add a reminder comment to the file by copying the line above and preceding it with `#`.) Each of the first four fields can be a literal value, or `*`, which matches anything. The password field from the first line that matches the current connection parameters will be used. (Therefore, put more-specific entries first when you are using wildcards.) If an entry needs to contain `:` or `\`, escape this character with `\`. The host name field is matched to the `host` connection parameter if that is specified, otherwise to the `hostaddr` parameter if that is specified; if neither are given then the host name `localhost` is searched for. The host name `localhost` is also searched for when the connection is a Unix-domain socket connection and the `host` parameter matches libpq's default socket directory path. In a standby server, a database field of `replication` matches streaming replication connections made to the master server. The database field is of limited usefulness otherwise, because users have the same password for all databases in the same cluster.

On Unix systems, the permissions on a password file must disallow any access to world or group; achieve this by a command such as `chmod 0600 ~/.pgpass`. If the permissions are less strict than this, the file will be ignored. On Microsoft Windows, it is assumed that the file is stored in a directory that is secure, so no special permissions check is made.

## 34.16. The Connection Service File

The connection service file allows libpq connection parameters to be associated with a single service name. That service name can then be specified by a libpq connection, and the associated settings will be used. This allows connection parameters to be modified without requiring a recompile of the libpq application. The service name can also be specified using the `PGSERVICE` environment variable.

The connection service file can be a per-user service file at `~/.pg_service.conf` or the location specified by the environment variable `PGSERVICEFILE`, or it can be a system-wide file at `~/.pgconfig --sysconfdir`/pg_service.conf or in the directory specified by the environment variable `PGSYSConfDIR`. If service definitions with the same name exist in the user and the system file, the user file takes precedence.

The file uses an “INI file” format where the section name is the service name and the parameters are connection parameters; see Section 34.1.2 for a list. For example:

```
# comment
[mydb]
host=somehost
port=5433
user=admin
```
An example file is provided at share/pg_service.conf.sample.

34.17. LDAP Lookup of Connection Parameters

If libpq has been compiled with LDAP support (option --with-ldap for configure) it is possible to retrieve connection options like host or dbname via LDAP from a central server. The advantage is that if the connection parameters for a database change, the connection information doesn’t have to be updated on all client machines.

LDAP connection parameter lookup uses the connection service file pg_service.conf (see Section 34.16). A line in a pg_service.conf stanza that starts with ldap:// will be recognized as an LDAP URL and an LDAP query will be performed. The result must be a list of keyword = value pairs which will be used to set connection options. The URL must conform to RFC 1959 and be of the form

ldap://[hostname[:port]]/search_base?attribute?search_scope?filter

where hostname defaults to localhost and port defaults to 389.

Processing of pg_service.conf is terminated after a successful LDAP lookup, but is continued if the LDAP server cannot be contacted. This is to provide a fallback with further LDAP URL lines that point to different LDAP servers, classical keyword = value pairs, or default connection options. If you would rather get an error message in this case, add a syntactically incorrect line after the LDAP URL.

A sample LDAP entry that has been created with the LDIF file

version:1
dn:cn=mydatabase,dc=mycompany,dc=com
changetype:add
objectclass:top
objectclass:device
cn:mydatabase
description:host=dbserver.mycompany.com
description:port=5439
description:dbname=mydb
description:user=mydb_user
description:sslmode=require

might be queried with the following LDAP URL:

ldap://ldap.mycompany.com/dc=mycompany,dc=com?description?one?
(cn=mydatabase)

You can also mix regular service file entries with LDAP lookups. A complete example for a stanza in pg_service.conf would be:

# only host and port are stored in LDAP, specify dbname and user explicitly
[customerdb]
dbname=customer
user=appuser
ldap://ldap.acme.com/cn=dbserver,cn=hosts?pgconnectinfo?base?
(objectclass=*)
34.18. SSL Support

PostgreSQL has native support for using SSL connections to encrypt client/server communications for increased security. See Section 18.9 for details about the server-side SSL functionality.

libpq reads the system-wide OpenSSL configuration file. By default, this file is named `openssl.cnf` and is located in the directory reported by `openssl version -d`. This default can be overridden by setting environment variable `OPENSSL_CONF` to the name of the desired configuration file.

34.18.1. Client Verification of Server Certificates

By default, PostgreSQL will not perform any verification of the server certificate. This means that it is possible to spoof the server identity (for example by modifying a DNS record or by taking over the server IP address) without the client knowing. In order to prevent spoofing, the client must be able to verify the server's identity via a chain of trust. A chain of trust is established by placing a root (self-signed) certificate authority (CA) certificate on one computer and a leaf certificate signed by the root certificate on another computer. It is also possible to use an “intermediate” certificate which is signed by the root certificate and signs leaf certificates.

To allow the client to verify the identity of the server, place a root certificate on the client and a leaf certificate signed by the root certificate on the server. To allow the server to verify the identity of the client, place a root certificate on the server and a leaf certificate signed by the root certificate on the client. One or more intermediate certificates (usually stored with the leaf certificate) can also be used to link the leaf certificate to the root certificate.

Once a chain of trust has been established, there are two ways for the client to validate the leaf certificate sent by the server. If the parameter `sslmode` is set to `verify-ca`, libpq will verify that the server is trustworthy by checking the certificate chain up to the root certificate stored on the client. If `sslmode` is set to `verify-full`, libpq will also verify that the server host name matches the name stored in the server certificate. The SSL connection will fail if the server certificate cannot be verified. `verify-full` is recommended in most security-sensitive environments.

In `verify-full` mode, the host name is matched against the certificate's Subject Alternative Name attribute(s), or against the Common Name attribute if no Subject Alternative Name of type dNSName is present. If the certificate's name attribute starts with an asterisk (*), the asterisk will be treated as a wildcard, which will match all characters except a dot (.). This means the certificate will not match subdomains. If the connection is made using an IP address instead of a host name, the IP address will be matched (without doing any DNS lookups).

To allow server certificate verification, one or more root certificates must be placed in the file `~/.postgresql/root.crt` in the user's home directory. (On Microsoft Windows the file is named `%APPDATA%\postgresql\root.crt`.) Intermediate certificates should also be added to the file if they are needed to link the certificate chain sent by the server to the root certificates stored on the client.

Certificate Revocation List (CRL) entries are also checked if the file `~/.postgresql/root.crl` exists (`%APPDATA%\postgresql\root.crl` on Microsoft Windows).

The location of the root certificate file and the CRL can be changed by setting the connection parameters `sslrootcert` and `sslcrl` or the environment variables `PGSSLROOTCERT` and `PGSSLCRL`.

Note

For backwards compatibility with earlier versions of PostgreSQL, if a root CA file exists, the behavior of `sslmode=require` will be the same as that of `verify-ca`, meaning the server certificate is validated against the CA. Relying on this behavior is discouraged, and applications that need certificate validation should always use `verify-ca` or `verify-full`. 
34.18.2. Client Certificates

If the server attempts to verify the identity of the client by requesting the client's leaf certificate, libpq will send the certificates stored in file ~/.postgresql/postgresql.crt in the user's home directory. The certificates must chain to the root certificate trusted by the server. A matching private key file ~/.postgresql/postgresql.key must also be present. The private key file must not allow any access to world or group; achieve this by the command chmod 0600 ~/.postgresql/postgresql.key. On Microsoft Windows these files are named %APPDATA%\postgresql\postgresql.crt and %APPDATA%\postgresql\postgresql.key, and there is no special permissions check since the directory is presumed secure. The location of the certificate and key files can be overridden by the connection parameters sslcert and sslkey or the environment variables PGSSLCERT and PGSSLKEY.

The first certificate in postgresql.crt must be the client's certificate because it must match the client's private key. "Intermediate" certificates can be optionally appended to the file — doing so avoids requiring storage of intermediate certificates on the server (ssl_ca_file).

For instructions on creating certificates, see Section 18.9.5.

34.18.3. Protection Provided in Different Modes

The different values for the sslmode parameter provide different levels of protection. SSL can provide protection against three types of attacks:

Eavesdropping

If a third party can examine the network traffic between the client and the server, it can read both connection information (including the user name and password) and the data that is passed. SSL uses encryption to prevent this.

Man in the middle (MITM)

If a third party can modify the data while passing between the client and server, it can pretend to be the server and therefore see and modify data even if it is encrypted. The third party can then forward the connection information and data to the original server, making it impossible to detect this attack. Common vectors to do this include DNS poisoning and address hijacking, whereby the client is directed to a different server than intended. There are also several other attack methods that can accomplish this. SSL uses certificate verification to prevent this, by authenticating the server to the client.

Impersonation

If a third party can pretend to be an authorized client, it can simply access data it should not have access to. Typically this can happen through insecure password management. SSL uses client certificates to prevent this, by making sure that only holders of valid certificates can access the server.

For a connection to be known secure, SSL usage must be configured on both the client and the server before the connection is made. If it is only configured on the server, the client may end up sending sensitive information (e.g. passwords) before it knows that the server requires high security. In libpq, secure connections can be ensured by setting the sslmode parameter to verify-full or verify-ca, and providing the system with a root certificate to verify against. This is analogous to using an https URL for encrypted web browsing.

Once the server has been authenticated, the client can pass sensitive data. This means that up until this point, the client does not need to know if certificates will be used for authentication, making it safe to specify that only in the server configuration.

All SSL options carry overhead in the form of encryption and key-exchange, so there is a trade-off that has to be made between performance and security. Table 34.1 illustrates the risks the different sslmode values protect against, and what statement they make about security and overhead.
Table 34.1. SSL Mode Descriptions

<table>
<thead>
<tr>
<th>sslmode</th>
<th>Eavesdropping protection</th>
<th>MITM protection</th>
<th>Statement</th>
</tr>
</thead>
<tbody>
<tr>
<td>disable</td>
<td>No</td>
<td>No</td>
<td>I don't care about security, and I don't want to pay the overhead of encryption.</td>
</tr>
<tr>
<td>allow</td>
<td>Maybe</td>
<td>No</td>
<td>I don't care about security, but I will pay the overhead of encryption if the server insists on it.</td>
</tr>
<tr>
<td>prefer</td>
<td>Maybe</td>
<td>No</td>
<td>I don't care about encryption, but I wish to pay the overhead of encryption if the server supports it.</td>
</tr>
<tr>
<td>require</td>
<td>Yes</td>
<td>No</td>
<td>I want my data to be encrypted, and I accept the overhead. I trust that the network will make sure I always connect to the server I want.</td>
</tr>
<tr>
<td>verify-ca</td>
<td>Yes</td>
<td>Depends on CA-policy</td>
<td>I want my data encrypted, and I accept the overhead. I want to be sure that I connect to a server that I trust.</td>
</tr>
<tr>
<td>verify-full</td>
<td>Yes</td>
<td>Yes</td>
<td>I want my data encrypted, and I accept the overhead. I want to be sure that I connect to a server I trust, and that it’s the one I specify.</td>
</tr>
</tbody>
</table>

The difference between verify-ca and verify-full depends on the policy of the root CA. If a public CA is used, verify-ca allows connections to a server that somebody else may have registered with the CA. In this case, verify-full should always be used. If a local CA is used, or even a self-signed certificate, using verify-ca often provides enough protection.

The default value for sslmode is prefer. As is shown in the table, this makes no sense from a security point of view, and it only promises performance overhead if possible. It is only provided as the default for backward compatibility, and is not recommended in secure deployments.

34.18.4. SSL Client File Usage

Table 34.2 summarizes the files that are relevant to the SSL setup on the client.

Table 34.2. Libpq/Client SSL File Usage

<table>
<thead>
<tr>
<th>File</th>
<th>Contents</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>~/.postgresql/postgresql.crt</td>
<td>client certificate</td>
<td>requested by server</td>
</tr>
<tr>
<td>~/.postgresql/postgresql.key</td>
<td>client private key</td>
<td>proves client certificate sent by owner; does not indicate certificate owner is trustworthy</td>
</tr>
</tbody>
</table>
### 34.18.5. SSL Library Initialization

If your application initializes libssl and/or libcrypto libraries and libpq is built with SSL support, you should call PQinitOpenSSL to tell libpq that the libssl and/or libcrypto libraries have been initialized by your application, so that libpq will not also initialize those libraries. See [http://h41379.www4.hpe.com/doc/83final/ba554_90007/ch04.html](http://h41379.www4.hpe.com/doc/83final/ba554_90007/ch04.html) for details on the SSL API.

```c
void PQinitOpenSSL(int do_ssl, int do_crypto);
```

This function allows applications to select which security libraries to initialize.

- **PQinitSSL**: This function is equivalent to `PQinitOpenSSL(do_ssl, do_ssl). It is sufficient for applications that initialize both or neither of OpenSSL and libcrypto.

- **PQinitSSL** has been present since PostgreSQL 8.0, while **PQinitOpenSSL** was added in PostgreSQL 8.4, so PQinitSSL might be preferable for applications that need to work with older versions of libpq.

### 34.19. Behavior in Threaded Programs

libpq is reentrant and thread-safe by default. You might need to use special compiler command-line options when you compile your application code. Refer to your system’s documentation for information about how to build thread-enabled applications, or look in `src/Makefile.global` for `PTHREAD_CFLAGS` and `PTHREAD_LIBS`. This function allows the querying of libpq’s thread-safe status:

```c
int PQisthreadsafe();
```
libpq - C Library

Returns 1 if the libpq is thread-safe and 0 if it is not.

One thread restriction is that no two threads attempt to manipulate the same PGconn object at the same time. In particular, you cannot issue concurrent commands from different threads through the same connection object. (If you need to run concurrent commands, use multiple connections.)

PGresult objects are normally read-only after creation, and so can be passed around freely between threads. However, if you use any of the PGresult-modifying functions described in Section 34.11 or Section 34.13, it's up to you to avoid concurrent operations on the same PGresult, too.

The deprecated functions PQrequestCancel and PQoidStatus are not thread-safe and should not be used in multithread programs. PQrequestCancel can be replaced by PQcancel. PQoidStatus can be replaced by PQoidValue.

If you are using Kerberos inside your application (in addition to inside libpq), you will need to do locking around Kerberos calls because Kerberos functions are not thread-safe. See function PQregisterThreadLock in the libpq source code for a way to do cooperative locking between libpq and your application.

If you experience problems with threaded applications, run the program in src/tools/thread to see if your platform has thread-unsafe functions. This program is run by configure, but for binary distributions your library might not match the library used to build the binaries.

34.20. Building libpq Programs

To build (i.e., compile and link) a program using libpq you need to do all of the following things:

• Include the libpq-fe.h header file:

#include <libpq-fe.h>

If you failed to do that then you will normally get error messages from your compiler similar to:

foo.c: In function `main':
foo.c:34: `PGconn' undeclared (first use in this function)
foo.c:35: `PGresult' undeclared (first use in this function)
foo.c:54: `CONNECTION_BAD' undeclared (first use in this function)
foo.c:68: `PGRES_COMMAND_OK' undeclared (first use in this function)
foo.c:95: `PGRES_TUPLES_OK' undeclared (first use in this function)

• Point your compiler to the directory where the PostgreSQL header files were installed, by supplying the -I directory option to your compiler. (In some cases the compiler will look into the directory in question by default, so you can omit this option.) For instance, your compile command line could look like:

cc -c -I/usr/local/pgsql/include testprog.c

If you are using makefiles then add the option to the CPPFLAGS variable:

CPPFLAGS += -I/usr/local/pgsql/include

If there is any chance that your program might be compiled by other users then you should not hardcode the directory location like that. Instead, you can run the utility pg_config to find out where the header files are on the local system:
$ pg_config --includedir
/usr/local/include

If you have pkg-config installed, you can run instead:

$ pkg-config --cflags libpq
-I/usr/local/include

Note that this will already include the -I in front of the path.

Failure to specify the correct option to the compiler will result in an error message such as:

testlibpq.c:8:22: libpq-fe.h: No such file or directory

• When linking the final program, specify the option `-lpq` so that the libpq library gets pulled in, as well as the option `-Ldirectory` to point the compiler to the directory where the libpq library resides. (Again, the compiler will search some directories by default.) For maximum portability, put the -L option before the -lpq option. For example:

cc -o testprog testprogl.o testprog2.o -L/usr/local/pgsql/lib -lpq

You can find out the library directory using pg_config as well:

$ pg_config --libdir
/usr/local/pgsql/lib

Or again use pkg-config:

$ pkg-config --libs libpq
-L/usr/local/pgsql/lib -lpq

Note again that this prints the full options, not only the path.

Error messages that point to problems in this area could look like the following:

```
testlibpq.o: In function `main':
testlibpq.o(.text+0x60): undefined reference to `PQsetdbLogin'
testlibpq.o(.text+0x71): undefined reference to `PQstatus'
testlibpq.o(.text+0xa4): undefined reference to `PQerrorMessage'
```

This means you forgot `-lpq`.

```
/usr/bin/ld: cannot find -lpq
```

This means you forgot the -L option or did not specify the right directory.

### 34.21. Example Programs

These examples and others can be found in the directory `src/test/examples` in the source code distribution.
Example 34.1. libpq Example Program 1

/*
 * src/test/examples/testlibpq.c
 *
 * testlibpq.c
 *
 * Test the C version of libpq, the PostgreSQL frontend library.
 */
#include <stdio.h>
#include <stdlib.h>
#include "libpq-fe.h"

static void
exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

int
main(int argc, char **argv)
{
    const char *conninfo;
    PGconn *conn;
    PGresult *res;
    int nFields;
    int i,
        j;

    /*
     * If the user supplies a parameter on the command line, use it
     * as the
     * conninfo string; otherwise default to setting
     * dbname=postgres and using
     * environment variables or defaults for all other connection
     * parameters.
     */
    if (argc > 1)
        conninfo = argv[1];
    else
        conninfo = "dbname = postgres";

    /* Make a connection to the database */
    conn = PQconnectdb(conninfo);

    /* Check to see that the backend connection was successfully
     * made */
    if (PQstatus(conn) != CONNECTION_OK)
    {
        fprintf(stderr, "Connection to database failed: %s",
                PQerrorMessage(conn));
        exit_nicely(conn);
    }
/* Set always-secure search path, so malicious users can't take control. */
res = PQexec(conn,
    "SELECT pg_catalog.set_config('search_path', '', false)");
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}

/*
* Should PQclear PGresult whenever it is no longer needed to avoid memory leaks
*/
PQclear(res);

/*
* Our test case here involves using a cursor, for which we must be inside a transaction block. We could do the whole thing with a single PQexec() of "select * from pg_database", but that's too trivial to make a good example.
*/

/* Start a transaction block */
res = PQexec(conn, "BEGIN");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "BEGIN command failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

/*
* Fetch rows from pg_database, the system catalog of databases
*/
res = PQexec(conn, "DECLARE myportal CURSOR FOR select * from pg_database");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "DECLARE CURSOR failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

res = PQexec(conn, "FETCH ALL in myportal");
if (PQresultStatus(res) != PGRES_TUPLES_OK)
Example 34.2. libpq Example Program 2

Example 34.2. libpq Example Program 2

/*
 * src/test/examples/testlibpq2.c
 * 
 * testlibpq2.c
 *   Test of the asynchronous notification interface
 * 
 * Start this program, then from psql in another window do
 *   NOTIFY TBL2;
 *   Repeat four times to get this program to exit.
 * 
 * Or, if you want to get fancy, try this:
 *   populate a database with the following commands
 *   (provided in src/test/examples/testlibpq2.sql):
 * 
 *   CREATE SCHEMA TESTLIBPQ2;
 */
libpq - C Library

* SET search_path = TESTLIBPQ2;
* CREATE TABLE TBL1 (i int4);
* CREATE TABLE TBL2 (i int4);
* CREATE RULE r1 AS ON INSERT TO TBL1 DO
  (INSERT INTO TBL2 VALUES (new.i); NOTIFY TBL2);
*
* Start this program, then from psql do this four times:
*
* INSERT INTO TESTLIBPQ2.TBL1 VALUES (10);
*/

#include WIN32
#include <windows.h>
#endif
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <errno.h>
#include <sys/time.h>
#include <sys/types.h>
#ifndef HAVE_SYS_SELECT_H
#include <sys/select.h>
#endif
#include "libpq-fe.h"

static void
exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

int
main(int argc, char **argv)
{
    const char *conninfo;
    PGconn    *conn;
    PGresult  *res;
    PGnotify  *notify;
    int       nnotifies;

    /*
     * If the user supplies a parameter on the command line, use it as the
     * conninfo string; otherwise default to setting
     * dbname=postgres and using
     * environment variables or defaults for all other connection
     * parameters.
     */
    if (argc > 1)
        conninfo = argv[1];
    else
        conninfo = "dbname = postgres";

    /* Make a connection to the database */
    conn = PQconnectdb(conninfo);
/ * Check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
            PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can't take control. */
res = PQexec(conn, "SELECT pg_catalog.set_config('search_path', '', false)");
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}

/* Should PQclear PGresult whenever it is no longer needed to avoid memory leaks */
PQclear(res);

/* Issue LISTEN command to enable notifications from the rule's NOTIFY. */
res = PQexec(conn, "LISTEN TBL2");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "LISTEN command failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

/* Quit after four notifies are received. */
nnotifies = 0;
while (nnotifies < 4)
{
    /*
    * Sleep until something happens on the connection. We use select(2) * to wait for input, but you could also use poll() or similar * facilities. *
    */
    int         sock;
    fd_set      input_mask;

    sock = PQsocket(conn);
if (sock < 0)
    break;              /* shouldn't happen */
FD_ZERO(&input_mask);
FD_SET(sock, &input_mask);
if (select(sock + 1, &input_mask, NULL, NULL, NULL) < 0)
{
    fprintf(stderr, "select() failed: %s\n",
            strerror(errno));
    exit_nicely(conn);
}
/* Now check for input */
PQconsumeInput(conn);
while ((notify = PQnotifies(conn)) != NULL)
{
    fprintf(stderr,
            "ASYNC NOTIFY of '\%s' received from backend PID
            \d\n",
            notify->relname, notify->be_pid);
    PQfreemem(notify);
    nnotifies++;
    PQconsumeInput(conn);
}

fprintf(stderr, "Done.\n");
/* close the connection to the database and cleanup */
PQfinish(conn);
return 0;
}

Example 34.3. libpq Example Program 3

/*
 * src/test/examples/testlibpq3.c
 *
 * testlibpq3.c
 * Test out-of-line parameters and binary I/O.
 *
 * Before running this, populate a database with the following commands
 * (provided in src/test/examples/testlibpq3.sql):
 * CREATE SCHEMA testlibpq3;
 * SET search_path = testlibpq3;
 * CREATE TABLE test1 (i int4, t text, b bytea);
 * INSERT INTO test1 values (1, 'joe''s place', '\000\001\002\003\004');
 * INSERT INTO test1 values (2, 'ho there', '\004\003\002\001\000');
 */
* The expected output is:
* tuple 0: got
  * i = (4 bytes) 1
  * t = (11 bytes) 'joe's place'
  * b = (5 bytes) \000\001\002\003\004
* tuple 0: got
  * i = (4 bytes) 2
  * t = (8 bytes) 'ho there'
  * b = (5 bytes) \004\003\002\001\000
*/

#ifdef WIN32
#include <windows.h>
#endif
#include <stdio.h>
#include <stdlib.h>
#include <stdint.h>
#include <string.h>
#include <sys/types.h>
#include "libpq-fe.h"
/* for ntohl/htonl */
#include <netinet/in.h>
#include <arpa/inet.h>

static void exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

/* This function prints a query result that is a binary-format fetch from
 * a table defined as in the comment above. We split it out because the
 * main() function uses it twice.
 */
static void show_binary_results(PGresult *res)
{
    int         i,
                j;
    int         i_fnum,
                t_fnum,
                b_fnum;

    /* Use PQfnumber to avoid assumptions about field order in result */
    i_fnum = PQfnumber(res, "i");
    t_fnum = PQfnumber(res, "t");
    b_fnum = PQfnumber(res, "b");
for (i = 0; i < PQntuples(res); i++)
{
    char       *iptr;
    char       *tptr;
    char       *bptr;
    int         blen;
    int         ival;

    /* Get the field values (we ignore possibility they are null!) */
    iptr = PQgetvalue(res, i, i_fnum);
    tptr = PQgetvalue(res, i, t_fnum);
    bptr = PQgetvalue(res, i, b_fnum);

    /*
     * The binary representation of INT4 is in network byte order, which
     * we'd better coerce to the local byte order.
     */
    ival = ntohl(*((uint32_t *) iptr));

    /*
     * The binary representation of TEXT is, well, text, and since libpq
     * was nice enough to append a zero byte to it, it'll work just fine
     * as a C string.
     *
     * The binary representation of BYTEA is a bunch of bytes, which could
     * include embedded nulls so we have to pay attention to field length.
     */
    blen = PQgetlength(res, i, b_fnum);

    printf("tuple %d: got\n", i);
    printf(" i = (%d bytes) %d\n", PQgetlength(res, i, i_fnum), ival);
    printf(" t = (%d bytes) '%s'\n", PQgetlength(res, i, t_fnum), tptr);
    printf(" b = (%d bytes) ", blen);
    for (j = 0; j < blen; j++)
        printf("\%03o", bptr[j]);
    printf("\n\n\n");
}

int
main(int argc, char **argv)
{
    const char *conninfo;
    PGconn     *conn;
    PGresult   *res;
    const char *paramValues[1];
    int         paramLengths[1];
    int         paramFormats[1];
    uint32_t    binaryIntVal;

    846
/*
 * If the user supplies a parameter on the command line, use it as the
 * conninfo string; otherwise default to setting dbname=postgres and using
 * environment variables or defaults for all other connection parameters.
 * /
if (argc > 1)
conninfo = argv[1];
else
conninfo = "dbname = postgres";

/* Make a connection to the database */
conn = PQconnectdb(conninfo);

/* Check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can't take control. */
res = PQexec(conn, "SET search_path = testlibpq3");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

/* The point of this program is to illustrate use of PQexecParams() with
 * out-of-line parameters, as well as binary transmission of data.
 *
 * This first example transmits the parameters as text, but receives the
 * results in binary format. By using out-of-line parameters we can avoid
 * a lot of tedious mucking about with quoting and escaping, even though
 * the data is text. Notice how we don't have to do anything special with
 * the quote mark in the parameter value.
 */

/* Here is our out-of-line parameter value */
paramValues[0] = "joe's place";
res = PQexecParams(conn,
"SELECT * FROM test1 WHERE t = $1",
1,          /* one param */
NULL,       /* let the backend deduce param
           */
paramValues,
NULL,       /* don't need param lengths since
           */
NULL,       /* default to all text params */
1);         /* ask for binary results */

if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SELECT failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}

show_binary_results(res);
PQclear(res);

/*
 * In this second example we transmit an integer parameter in
 * binary form,
 * and again retrieve the results in binary form.
 * Although we tell PQexecParams we are letting the backend
deduce
 * parameter type, we really force the decision by casting the
 * parameter
 * symbol in the query text. This is a good safety measure
 * when sending
 * binary parameters.
 */

/* Convert integer value "2" to network byte order */
binaryIntVal = htonl((uint32_t) 2);

/* Set up parameter arrays for PQexecParams */
paramValues[0] = (char *) &binaryIntVal;
paramLengths[0] = sizeof(binaryIntVal);
paramFormats[0] = 1;        /* binary */
res = PQexecParams(conn,
    "SELECT * FROM test1 WHERE i = $1::int4",
    1,       /* one param */
    NULL,    /* let the backend deduce param
           */
    paramValues,
    paramLengths,
    paramFormats,
    1);      /* ask for binary results */

if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SELECT failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
show_binary_results(res);
PQclear(res);

/* close the connection to the database and cleanup */
PQfinish(conn);

return 0;
Chapter 35. Large Objects

PostgreSQL has a *large object* facility, which provides stream-style access to user data that is stored in a special large-object structure. Streaming access is useful when working with data values that are too large to manipulate conveniently as a whole.

This chapter describes the implementation and the programming and query language interfaces to PostgreSQL large object data. We use the libpq C library for the examples in this chapter, but most programming interfaces native to PostgreSQL support equivalent functionality. Other interfaces might use the large object interface internally to provide generic support for large values. This is not described here.

### 35.1. Introduction

All large objects are stored in a single system table named `pg_largeobject`. Each large object also has an entry in the system table `pg_largeobject_metadata`. Large objects can be created, modified, and deleted using a read/write API that is similar to standard operations on files.

PostgreSQL also supports a storage system called “TOAST”, which automatically stores values larger than a single database page into a secondary storage area per table. This makes the large object facility partially obsolete. One remaining advantage of the large object facility is that it allows values up to 4 TB in size, whereas TOASTed fields can be at most 1 GB. Also, reading and updating portions of a large object can be done efficiently, while most operations on a TOASTed field will read or write the whole value as a unit.

### 35.2. Implementation Features

The large object implementation breaks large objects up into “chunks” and stores the chunks in rows in the database. A B-tree index guarantees fast searches for the correct chunk number when doing random access reads and writes.

The chunks stored for a large object do not have to be contiguous. For example, if an application opens a new large object, seeks to offset 1000000, and writes a few bytes there, this does not result in allocation of 1000000 bytes worth of storage; only of chunks covering the range of data bytes actually written. A read operation will, however, read out zeroes for any unallocated locations preceding the last existing chunk. This corresponds to the common behavior of “sparsely allocated” files in Unix file systems.

As of PostgreSQL 9.0, large objects have an owner and a set of access permissions, which can be managed using GRANT and REVOKE. SELECT privileges are required to read a large object, and UPDATE privileges are required to write or truncate it. Only the large object’s owner (or a database superuser) can delete, comment on, or change the owner of a large object. To adjust this behavior for compatibility with prior releases, see the `lo_compat_privileges` run-time parameter.

### 35.3. Client Interfaces

This section describes the facilities that PostgreSQL’s libpq client interface library provides for accessing large objects. The PostgreSQL large object interface is modeled after the Unix file-system interface, with analogues of `open`, `read`, `write`, `lseek`, etc.

All large object manipulation using these functions *must* take place within an SQL transaction block, since large object file descriptors are only valid for the duration of a transaction.

If an error occurs while executing any one of these functions, the function will return an otherwise-impossible value, typically 0 or -1. A message describing the error is stored in the connection object and can be retrieved with `PQerrorMessage`.

850
Client applications that use these functions should include the header file libpq/libpq-fs.h and link with the libpq library.

### 35.3.1. Creating a Large Object

The function

\[
\text{Oid lo\_creat(PGconn *conn, int mode);}\
\]

creates a new large object. The return value is the OID that was assigned to the new large object, or InvalidOid (zero) on failure. \text{mode} is unused and ignored as of PostgreSQL 8.1; however, for backward compatibility with earlier releases it is best to set it to INV\_READ, INV\_WRITE, or INV\_READ | INV\_WRITE. (These symbolic constants are defined in the header file libpq/libpq-fs.h.)

An example:

\[
\text{inv\_oid = lo\_creat(conn, INV\_READ|INV\_WRITE);}\
\]

The function

\[
\text{Oid lo\_create(PGconn *conn, Oid lobjId);}\
\]

also creates a new large object. The OID to be assigned can be specified by \text{lobjId}; if so, failure occurs if that OID is already in use for some large object. If \text{lobjId} is \text{InvalidOid} (zero) then \text{lo\_create} assigns an unused OID (this is the same behavior as \text{lo\_creat}). The return value is the OID that was assigned to the new large object, or \text{InvalidOid} (zero) on failure.

\text{lo\_create} is new as of PostgreSQL 8.1; if this function is run against an older server version, it will fail and return \text{InvalidOid}.

An example:

\[
\text{inv\_oid = lo\_create(conn, desired\_oid);}\
\]

### 35.3.2. Importing a Large Object

To import an operating system file as a large object, call

\[
\text{Oid lo\_import(PGconn *conn, const char *filename);}\
\]

\text{filename} specifies the operating system name of the file to be imported as a large object. The return value is the OID that was assigned to the new large object, or \text{InvalidOid} (zero) on failure. Note that the file is read by the client interface library, not by the server; so it must exist in the client file system and be readable by the client application.

The function

\[
\text{Oid lo\_import\_with\_oid(PGconn *conn, const char *filename, Oid lobjId);}\
\]

also imports a new large object. The OID to be assigned can be specified by \text{lobjId}; if so, failure occurs if that OID is already in use for some large object. If \text{lobjId} is \text{InvalidOid} (zero) then \text{lo\_import\_with\_oid} assigns an unused OID (this is the same behavior as \text{lo\_import}). The return value is the OID that was assigned to the new large object, or \text{InvalidOid} (zero) on failure.
lo_import_with_oid is new as of PostgreSQL 8.4 and uses lo_create internally which is new in 8.1; if this function is run against 8.0 or before, it will fail and return InvalidOid.

35.3.3. Exporting a Large Object

To export a large object into an operating system file, call

```c
int lo_export(PGconn *conn, Oid lobjId, const char *filename);
```

The lobjId argument specifies the OID of the large object to export and the filename argument specifies the operating system name of the file. Note that the file is written by the client interface library, not by the server. Returns 1 on success, -1 on failure.

35.3.4. Opening an Existing Large Object

To open an existing large object for reading or writing, call

```c
int lo_open(PGconn *conn, Oid lobjId, int mode);
```

The lobjId argument specifies the OID of the large object to open. The mode bits control whether the object is opened for reading (INV_READ), writing (INV_WRITE), or both. (These symbolic constants are defined in the header file libpq/libpq-fs.h.) lo_open returns a (non-negative) large object descriptor for later use in lo_read, lo_write, lo_lseek, lo_lseek64, lo_tell, lo_tell64, lo_truncate, lo_truncate64, and lo_close. The descriptor is only valid for the duration of the current transaction. On failure, -1 is returned.

The server currently does not distinguish between modes INV_WRITE and INV_READ | INV_WRITE: you are allowed to read from the descriptor in either case. However there is a significant difference between these modes and INV_READ alone: with INV_READ you cannot write on the descriptor, and the data read from it will reflect the contents of the large object at the time of the transaction snapshot that was active when lo_open was executed, regardless of later writes by this or other transactions. Reading from a descriptor opened with INV_WRITE returns data that reflects all writes of other committed transactions as well as writes of the current transaction. This is similar to the behavior of REPEATABLE READ versus READ COMMITTED transaction modes for ordinary SQL SELECT commands.

lo_open will fail if SELECT privilege is not available for the large object, or if INV_WRITE is specified and UPDATE privilege is not available. (Prior to PostgreSQL 11, these privilege checks were instead performed at the first actual read or write call using the descriptor.) These privilege checks can be disabled with the lo_compat_privileges run-time parameter.

An example:

```c
inv_fd = lo_open(conn, inv_oid, INV_READ|INV_WRITE);
```

35.3.5. Writing Data to a Large Object

The function

```c
int lo_write(PGconn *conn, int fd, const char *buf, size_t len);
```

writes len bytes from buf (which must be of size len) to large object descriptor fd. The fd argument must have been returned by a previous lo_open. The number of bytes actually written is returned (in the current implementation, this will always equal len unless there is an error). In the event of an error, the return value is -1.
Although the `len` parameter is declared as `size_t`, this function will reject length values larger than `INT_MAX`. In practice, it's best to transfer data in chunks of at most a few megabytes anyway.

### 35.3.6. Reading Data from a Large Object

The function

```c
int lo_read(PGconn *conn, int fd, char *buf, size_t len);
```

reads up to `len` bytes from large object descriptor `fd` into `buf` (which must be of size `len`). The `fd` argument must have been returned by a previous `lo_open`. The number of bytes actually read is returned; this will be less than `len` if the end of the large object is reached first. In the event of an error, the return value is `-1`.

Although the `len` parameter is declared as `size_t`, this function will reject length values larger than `INT_MAX`. In practice, it's best to transfer data in chunks of at most a few megabytes anyway.

### 35.3.7. Seeking in a Large Object

To change the current read or write location associated with a large object descriptor, call

```c
int lo_lseek(PGconn *conn, int fd, int offset, int whence);
```

This function moves the current location pointer for the large object descriptor identified by `fd` to the new location specified by `offset`. The valid values for `whence` are `SEEK_SET` (seek from object start), `SEEK_CUR` (seek from current position), and `SEEK_END` (seek from object end). The return value is the new location pointer, or `-1` on error.

When dealing with large objects that might exceed 2GB in size, instead use

```c
pg_int64 lo_lseek64(PGconn *conn, int fd, pg_int64 offset, int whence);
```

This function has the same behavior as `lo_lseek`, but it can accept an `offset` larger than 2GB and/or deliver a result larger than 2GB. Note that `lo_lseek` will fail if the new location pointer would be greater than 2GB.

`lo_lseek64` is new as of PostgreSQL 9.3. If this function is run against an older server version, it will fail and return `-1`.

### 35.3.8. Obtaining the Seek Position of a Large Object

To obtain the current read or write location of a large object descriptor, call

```c
int lo_tell(PGconn *conn, int fd);
```

If there is an error, the return value is `-1`.

When dealing with large objects that might exceed 2GB in size, instead use

```c
pg_int64 lo_tell64(PGconn *conn, int fd);
```

This function has the same behavior as `lo_tell`, but it can deliver a result larger than 2GB. Note that `lo_tell` will fail if the current read/write location is greater than 2GB.

`lo_tell64` is new as of PostgreSQL 9.3. If this function is run against an older server version, it will fail and return `-1`. 
35.3.9. Truncating a Large Object

To truncate a large object to a given length, call

```c
int lo_truncate(PGconn *conn, int fd, size_t len);
```

This function truncates the large object descriptor `fd` to length `len`. The `fd` argument must have been returned by a previous `lo_open`. If `len` is greater than the large object's current length, the large object is extended to the specified length with null bytes (`\0`). On success, `lo_truncate` returns zero. On error, the return value is -1.

The read/write location associated with the descriptor `fd` is not changed.

Although the `len` parameter is declared as `size_t`, `lo_truncate` will reject length values larger than `INT_MAX`.

When dealing with large objects that might exceed 2GB in size, instead use

```c
int lo_truncate64(PGconn *conn, int fd, pg_int64 len);
```

This function has the same behavior as `lo_truncate`, but it can accept a `len` value exceeding 2GB.

`lo_truncate` is new as of PostgreSQL 8.3; if this function is run against an older server version, it will fail and return -1.

`lo_truncate64` is new as of PostgreSQL 9.3; if this function is run against an older server version, it will fail and return -1.

35.3.10. Closing a Large Object Descriptor

A large object descriptor can be closed by calling

```c
int lo_close(PGconn *conn, int fd);
```

where `fd` is a large object descriptor returned by `lo_open`. On success, `lo_close` returns zero. On error, the return value is -1.

Any large object descriptors that remain open at the end of a transaction will be closed automatically.

35.3.11. Removing a Large Object

To remove a large object from the database, call

```c
int lo_unlink(PGconn *conn, Oid lobjId);
```

The `lobjId` argument specifies the OID of the large object to remove. Returns 1 if successful, -1 on failure.

35.4. Server-side Functions

Server-side functions tailored for manipulating large objects from SQL are listed in Table 35.1.

**Table 35.1. SQL-oriented Large Object Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>lo_from_oid(lo_id)</code></td>
<td>oid</td>
<td>Create a large object and store data</td>
<td><code>\x528 \x84528</code></td>
<td><code>0x528\x84528</code></td>
</tr>
<tr>
<td><code>lo_from_bytea(lo_id)</code></td>
<td></td>
<td>Create a large object and store data</td>
<td><code>\x00\xff\xff\xff\xff</code></td>
<td><code>0x00\xff\xff\xff\xff</code></td>
</tr>
</tbody>
</table>
### Large Objects

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>oid, string</code></td>
<td><code>bytea</code></td>
<td>ta there, returning its OID. Pass 0 to have the system choose an OID.</td>
<td></td>
</tr>
<tr>
<td><code>lo_put</code></td>
<td><code>void</code></td>
<td>Write data at the given offset.</td>
<td><code>lo_put(24528, 1, '\xa')</code></td>
</tr>
<tr>
<td><code>oid, offset</code></td>
<td><code>bigint</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>str</code></td>
<td><code>bytea</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_get</code></td>
<td><code>bytea</code></td>
<td>Extract contents or a substring thereof.</td>
<td>`lo_get(24528, \xffaaff)</td>
</tr>
<tr>
<td><code>oid []</code></td>
<td><code>from</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>for</code></td>
<td><code>int</code></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_creat</code></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_create</code></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_unlink</code></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_import</code></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>lo_export</code></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

There are additional server-side functions corresponding to each of the client-side functions described earlier; indeed, for the most part the client-side functions are simply interfaces to the equivalent server-side functions. The ones just as convenient to call via SQL commands are `lo_creat`, `lo_create`, `lo_unlink`, `lo_import`, and `lo_export`. Here are examples of their use:

```sql
CREATE TABLE image (
    name            text,
    raster          oid
);

SELECT lo_creat(-1);       -- returns OID of new, empty large object

SELECT lo_create(43213);   -- attempts to create large object with OID 43213

SELECT lo_unlink(173454);  -- deletes large object with OID 173454

INSERT INTO image (name, raster)
VALUES ('beautiful image', lo_import('/etc/motd'));

INSERT INTO image (name, raster)  -- same as above, but specify OID to use
VALUES ('beautiful image', lo_import('/etc/motd', 68583));

SELECT lo_export(image.raster, '/tmp/motd') FROM image
WHERE name = 'beautiful image';
```

The server-side `lo_import` and `lo_export` functions behave considerably differently from their client-side analogs. These two functions read and write files in the server’s file system, using the permissions of the database’s owning user. Therefore, by default their use is restricted to superusers. In contrast, the client-side import and export functions read and write files in the client’s file system, using the permissions of the client program. The client-side functions do not require any database privileges, except the privilege to read or write the large object in question.

---

**Caution**

It is possible to GRANT use of the server-side `lo_import` and `lo_export` functions to non-superusers, but careful consideration of the security implications is required. A malicious user of such privileges could easily parlay them into becoming superuser (for example by rewriting server configuration files), or could attack the rest of the server’s file system without bothering to obtain database superuser privileges.
as such. Access to roles having such privilege must therefore be guarded just as carefully as access to superuser roles. Nonetheless, if use of server-side `lo_import` or `lo_export` is needed for some routine task, it’s safer to use a role with such privileges than one with full superuser privileges, as that helps to reduce the risk of damage from accidental errors.

The functionality of `lo_read` and `lo_write` is also available via server-side calls, but the names of the server-side functions differ from the client side interfaces in that they do not contain underscores. You must call these functions as `loread` and `lowrite`.

### 35.5. Example Program

Example 35.1 is a sample program which shows how the large object interface in libpq can be used. Parts of the program are commented out but are left in the source for the reader’s benefit. This program can also be found in `src/test/examples/testlo.c` in the source distribution.

**Example 35.1. Large Objects with libpq Example Program**

```c
#include <stdio.h>
#include <stdlib.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <unistd.h>
#include "libpq-fe.h"
#include "libpq/libpq-fs.h"

#define BUFSIZE 1024

importFile(PGconn *conn, char *filename)
```
Large Objects

```c

Oid lobjId;
int lobj_fd;
char buf[BUFSIZE];
int nbytes,
tmp;
int fd;

/*
 * open the file to be read in
 */
fd = open(filename, O_RDONLY, 0666);
if (fd < 0)
{
    /* error */
    fprintf(stderr, "cannot open unix file\"%s\"\n", filename);
}

/*
 * create the large object
 */
lobjId = lo_creat(conn, INV_READ | INV_WRITE);
if (lobjId == 0)
    fprintf(stderr, "cannot create large object");

lobj_fd = lo_open(conn, lobjId, INV_WRITE);

/*
 * read in from the Unix file and write to the inversion file
 */
while ((nbytes = read(fd, buf, BUFSIZE)) > 0)
{
    tmp = lo_write(conn, lobj_fd, buf, nbytes);
    if (tmp < nbytes)
        fprintf(stderr, "error while reading \"%s\", filename);
}

close(fd);
lo_close(conn, lobj_fd);

return lobjId;
}

static void
pickout(PGconn *conn, Oid lobjId, int start, int len)
{
    int lobj_fd;
    char *buf;
    int nbytes;
    int nread;

    lobj_fd = lo_open(conn, lobjId, INV_READ);
    if (lobj_fd < 0)
        fprintf(stderr, "cannot open large object %u", lobjId);

    lo_lseek(conn, lobj_fd, start, SEEK_SET);
    buf = malloc(len + 1);
```
nread = 0;
while (len - nread > 0)
{
    nbytes = lo_read(conn, lobj_fd, buf, len - nread);
    buf[nbytes] = '\0';
    fprintf(stderr, ">>> %s", buf);
    nread += nbytes;
    if (nbytes <= 0)
        break;               /* no more data? */
}
free(buf);
fprintf(stderr, "\n");
lo_close(conn, lobj_fd);
}

static void overwrite(PGconn *conn, Oid lobjId, int start, int len)
{
    int lobj_fd;
    char *buf;
    int nbytes;
    int nwritten;
    int i;

    lobj_fd = lo_open(conn, lobjId, INV_WRITE);
    if (lobj_fd < 0)
        fprintf(stderr, "cannot open large object %u", lobjId);
    lo_lseek(conn, lobj_fd, start, SEEK_SET);
    buf = malloc(len + 1);
    for (i = 0; i < len; i++)
        buf[i] = 'X';
    buf[i] = '\0';

    nwritten = 0;
    while (len - nwritten > 0)
    {
        nbytes = lo_write(conn, lobj_fd, buf + nwritten, len -
                          nwritten);
        nwritten += nbytes;
        if (nbytes <= 0)
        {
            fprintf(stderr, "\nWRITE FAILED!\n");
            break;
        }
    }
    free(buf);
    fprintf(stderr, "\n");
    lo_close(conn, lobj_fd);
}

/* exportFile -
 * export large object "lobjOid" to file "out_filename"
 */
static void
exportFile(PGconn *conn, Oid lobjId, char *filename)
{
    int         lobj_fd;
    char        buf[BUFSIZE];
    int         nbytes, tmp;
    int         fd;

    /*
    * open the large object
    */
    lobj_fd = lo_open(conn, lobjId, INV_READ);
    if (lobj_fd < 0)
        fprintf(stderr, "cannot open large object \%lu", lobjId);

    /*
    * open the file to be written to
    */
    fd = open(filename, O_CREAT | O_WRONLY | O_TRUNC, 0666);
    if (fd < 0) {                           /* error */
        /* error */
        fprintf(stderr, "cannot open unix file\"%s\"", filename);
    }

    /*
    * read in from the inversion file and write to the Unix file
    */
    while ((nbytes = lo_read(conn, lobj_fd, buf, BUFSIZE)) > 0)
    {
        tmp = write(fd, buf, nbytes);
        if (tmp < nbytes)
        {
            fprintf(stderr, "error while writing \"%s\", filename);
        }
    }

    lo_close(conn, lobj_fd);
    close(fd);

    return;
}

static void
exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

int
main(int argc, char **argv)
{
    char     *in_filename,
              *out_filename;
    char     *database;

Oid lobjOid;
PGconn *conn;
PGresult *res;

if (argc != 4)
{
    fprintf(stderr, "Usage: %s database_name in_filename
out_filename\n", argv[0]);
    exit(1);
}

database = argv[1];
in_filename = argv[2];
out_filename = argv[3];

/* set up the connection */
conn = PQsetdb(NULL, NULL, NULL, NULL, database);

/* check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can't take control. */
res = PQexec(conn, "SELECT pg_catalog.set_config('search_path', '',
false));
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

res = PQexec(conn, "begin");
PQclear(res);
printf("importing file \"%s\" ...
", in_filename);
/* lobjOid = importFile(conn, in_filename); */
lobjOid = lo_import(conn, in_filename);
if (lobjOid == 0)
    fprintf(stderr, "%s\n", PQerrorMessage(conn));
else
{
    printf("there large object %u.\n", lobjOid);

    printf("picking out bytes 1000-2000 of the large object\n");
pickout(conn, lobjOid, 1000, 1000);
printf("overwriting bytes 1000-2000 of the large object
with X's\n");
overwrite(conn, lobjOid, 1000, 1000);

printf("exporting large object to file \"%s\" ...
", out_filename);
/*      exportFile(conn, lobjOid, out_filename); */
if (lo_export(conn, lobjOid, out_filename) < 0)
    fprintf(stderr, "%s\n", PQerrorMessage(conn));
}

res = PQexec(conn, "end");
PQclear(res);
PQfinish(conn);
return 0;
}
Chapter 36. ECPG - Embedded SQL in C

This chapter describes the embedded SQL package for PostgreSQL. It was written by Linus Tolke (<linus@epact.se>) and Michael Meskes (<meskes@postgresql.org>). Originally it was written to work with C. It also works with C++, but it does not recognize all C++ constructs yet.

This documentation is quite incomplete. But since this interface is standardized, additional information can be found in many resources about SQL.

36.1. The Concept

An embedded SQL program consists of code written in an ordinary programming language, in this case C, mixed with SQL commands in specially marked sections. To build the program, the source code (*.pgc) is first passed through the embedded SQL preprocessor, which converts it to an ordinary C program (*.c), and afterwards it can be processed by a C compiler. (For details about the compiling and linking see Section 36.10). Converted ECPG applications call functions in the libpq library through the embedded SQL library (ecpglib), and communicate with the PostgreSQL server using the normal frontend-backend protocol.

Embedded SQL has advantages over other methods for handling SQL commands from C code. First, it takes care of the tedious passing of information to and from variables in your C program. Second, the SQL code in the program is checked at build time for syntactical correctness. Third, embedded SQL in C is specified in the SQL standard and supported by many other SQL database systems. The PostgreSQL implementation is designed to match this standard as much as possible, and it is usually possible to port embedded SQL programs written for other SQL databases to PostgreSQL with relative ease.

As already stated, programs written for the embedded SQL interface are normal C programs with special code inserted to perform database-related actions. This special code always has the form:

EXEC SQL ...;

These statements syntactically take the place of a C statement. Depending on the particular statement, they can appear at the global level or within a function. Embedded SQL statements follow the case-sensitivity rules of normal SQL code, and not those of C. Also they allow nested C-style comments that are part of the SQL standard. The C part of the program, however, follows the C standard of not accepting nested comments.

The following sections explain all the embedded SQL statements.

36.2. Managing Database Connections

This section describes how to open, close, and switch database connections.

36.2.1. Connecting to the Database Server

One connects to a database using the following statement:

EXEC SQL CONNECT TO target [AS connection-name] [USER user-name];

The target can be specified in the following ways:
• `dbname[@hostname][:port]`
• `tcp:postgresql://hostname[:port][/dbname][?options]`
• `unix:postgresql://hostname[:port][/dbname][?options]`
• an SQL string literal containing one of the above forms
• a reference to a character variable containing one of the above forms (see examples)
• DEFAULT

If you specify the connection target literally (that is, not through a variable reference) and you don’t quote the value, then the case-insensitivity rules of normal SQL are applied. In that case you can also double-quote the individual parameters separately as needed. In practice, it is probably less error-prone to use a (single-quoted) string literal or a variable reference. The connection target DEFAULT initiates a connection to the default database under the default user name. No separate user name or connection name can be specified in that case.

There are also different ways to specify the user name:

• `username`
• `username/password`
• `username IDENTIFIED BY password`
• `username USING password`

As above, the parameters `username` and `password` can be an SQL identifier, an SQL string literal, or a reference to a character variable.

The `connection-name` is used to handle multiple connections in one program. It can be omitted if a program uses only one connection. The most recently opened connection becomes the current connection, which is used by default when an SQL statement is to be executed (see later in this chapter).

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from `search_path`. For example, add `options=-csearch_path= to options, or issue EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); after connecting. This consideration is not specific to ECPG; it applies to every interface for executing arbitrary SQL commands.

Here are some examples of CONNECT statements:

```sql
EXEC SQL CONNECT TO mydb@sql.mydomain.com;
EXEC SQL CONNECT TO unix:postgresql://sql.mydomain.com/mydb AS myconnection USER john;
EXEC SQL BEGIN DECLARE SECTION;
const char * target = "mydb@sql.mydomain.com";
const char * user = "john";
const char * passwd = "secret";
EXEC SQL END DECLARE SECTION;
```

```c
EXEC SQL CONNECT TO :target USER :user USING :passwd;
/* or EXEC SQL CONNECT TO :target USER :user/:passwd; */
```

The last form makes use of the variant referred to above as character variable reference. You will see in later sections how C variables can be used in SQL statements when you prefix them with a colon.
Be advised that the format of the connection target is not specified in the SQL standard. So if you want to develop portable applications, you might want to use something based on the last example above to encapsulate the connection target string somewhere.

### 36.2.2. Choosing a Connection

SQL statements in embedded SQL programs are by default executed on the current connection, that is, the most recently opened one. If an application needs to manage multiple connections, then there are two ways to handle this.

The first option is to explicitly choose a connection for each SQL statement, for example:

```sql
EXEC SQL AT connection-name SELECT ...;
```

This option is particularly suitable if the application needs to use several connections in mixed order.

If your application uses multiple threads of execution, they cannot share a connection concurrently. You must either explicitly control access to the connection (using mutexes) or use a connection for each thread.

The second option is to execute a statement to switch the current connection. That statement is:

```sql
EXEC SQL SET CONNECTION connection-name;
```

This option is particularly convenient if many statements are to be executed on the same connection.

Here is an example program managing multiple database connections:

```c
#include <stdio.h>
EXEC SQL BEGIN DECLARE SECTION;
char dbname[1024];
EXEC SQL END DECLARE SECTION;

int main()
{
    EXEC SQL CONNECT TO testdb1 AS con1 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
    EXEC SQL CONNECT TO testdb2 AS con2 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
    EXEC SQL CONNECT TO testdb3 AS con3 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

    /* This query would be executed in the last opened database "testdb3". */
    EXEC SQL SELECT current_database() INTO :dbname;
    printf("current=%.s (should be testdb3)\n", dbname);

    /* Using "AT" to run a query in "testdb2" */
    EXEC SQL AT con2 SELECT current_database() INTO :dbname;
    printf("current=%.s (should be testdb2)\n", dbname);

    return 0;
}
```

864
/* Switch the current connection to "testdb1". */
EXEC SQL SET CONNECTION con1;

EXEC SQL SELECT current_database() INTO :dbname;
printf("current=%s (should be testdb1)\n", dbname);

EXEC SQL DISCONNECT ALL;
return 0;
}

This example would produce this output:

current=testdb3  (should be testdb3)
current=testdb2  (should be testdb2)
current=testdb1  (should be testdb1)

### 36.2.3. Closing a Connection

To close a connection, use the following statement:

```sql
EXEC SQL DISCONNECT [connection];
```

The `connection` can be specified in the following ways:

- `connection-name`
- `DEFAULT`
- `CURRENT`
- `ALL`

If no connection name is specified, the current connection is closed.

It is good style that an application always explicitly disconnect from every connection it opened.

### 36.3. Running SQL Commands

Any SQL command can be run from within an embedded SQL application. Below are some examples of how to do that.

#### 36.3.1. Executing SQL Statements

Creating a table:

```sql
EXEC SQL CREATE TABLE foo (number integer, ascii char(16));
EXEC SQL CREATE UNIQUE INDEX num1 ON foo(number);
EXEC SQL COMMIT;
```

Inserting rows:

```sql
EXEC SQL INSERT INTO foo (number, ascii) VALUES (9999, 'doodad');
EXEC SQL COMMIT;
```

Deleting rows:
EXEC SQL DELETE FROM foo WHERE number = 9999;
EXEC SQL COMMIT;

Updates:

EXEC SQL UPDATE foo
    SET ascii = 'foobar'
    WHERE number = 9999;
EXEC SQL COMMIT;

SELECT statements that return a single result row can also be executed using EXEC SQL directly.
To handle result sets with multiple rows, an application has to use a cursor; see Section 36.3.2 below.
(As a special case, an application can fetch multiple rows at once into an array host variable; see Section 36.4.4.3.1.)

Single-row select:

EXEC SQL SELECT foo INTO :FooBar FROM table1 WHERE ascii = 'doodad';

Also, a configuration parameter can be retrieved with the SHOW command:

EXEC SQL SHOW search_path INTO :var;

The tokens of the form :something are host variables, that is, they refer to variables in the C program. They are explained in Section 36.4.

36.3.2. Using Cursors

To retrieve a result set holding multiple rows, an application has to declare a cursor and fetch each row from the cursor. The steps to use a cursor are the following: declare a cursor, open it, fetch a row from the cursor, repeat, and finally close it.

Select using cursors:

EXEC SQL DECLARE foo_bar CURSOR FOR
    SELECT number, ascii FROM foo
    ORDER BY ascii;
EXEC SQL OPEN foo_bar;
EXEC SQL FETCH foo_bar INTO :FooBar, DooDad;
...
EXEC SQL CLOSE foo_bar;
EXEC SQL COMMIT;

For more details about declaration of the cursor, see DECLARE, and see FETCH for FETCH command details.

Note

The ECPG DECLARE command does not actually cause a statement to be sent to the PostgreSQL backend. The cursor is opened in the backend (using the backend’s DECLARE command) at the point when the OPEN command is executed.
36.3.3. Managing Transactions

In the default mode, statements are committed only when EXEC SQL COMMIT is issued. The embedded SQL interface also supports autocommit of transactions (similar to psql’s default behavior) via the -t command-line option to ecpg (see ecpg) or via the EXEC SQL SET AUTOCOMMIT TO ON statement. In autocommit mode, each command is automatically committed unless it is inside an explicit transaction block. This mode can be explicitly turned off using EXEC SQL SET AUTOCOMMIT TO OFF.

The following transaction management commands are available:

EXEC SQL COMMIT
   Commit an in-progress transaction.

EXEC SQL ROLLBACK
   Roll back an in-progress transaction.

EXEC SQL PREPARE TRANSACTION transaction_id
   Prepare the current transaction for two-phase commit.

EXEC SQL COMMIT PREPARED transaction_id
   Commit a transaction that is in prepared state.

EXEC SQL ROLLBACK PREPARED transaction_id
   Roll back a transaction that is in prepared state.

EXEC SQL SET AUTOCOMMIT TO ON
   Enable autocommit mode.

EXEC SQL SET AUTOCOMMIT TO OFF
   Disable autocommit mode. This is the default.

36.3.4. Prepared Statements

When the values to be passed to an SQL statement are not known at compile time, or the same statement is going to be used many times, then prepared statements can be useful.

The statement is prepared using the command PREPARE. For the values that are not known yet, use the placeholder “?”:

EXEC SQL PREPARE stmt1 FROM "SELECT oid, datname FROM pg_database WHERE oid = ?";

If a statement returns a single row, the application can call EXECUTE after PREPARE to execute the statement, supplying the actual values for the placeholders with a USING clause:

EXEC SQL EXECUTE stmt1 INTO :dboid, :dbname USING 1;

If a statement returns multiple rows, the application can use a cursor declared based on the prepared statement. To bind input parameters, the cursor must be opened with a USING clause:

EXEC SQL PREPARE stmt1 FROM "SELECT oid, datname FROM pg_database WHERE oid > ?";
EXEC SQL DECLARE foo_bar CURSOR FOR stmt1;

/* when end of result set reached, break out of while loop */
EXEC SQL WHENEVER NOT FOUND DO BREAK;

EXEC SQL OPEN foo_bar USING 100;
...
while (1)
{
    EXEC SQL FETCH NEXT FROM foo_bar INTO :dboid, :dbname;
    ...
}
EXEC SQL CLOSE foo_bar;

When you don't need the prepared statement anymore, you should deallocate it:

EXEC SQL DEALLOCATE PREPARE name;

For more details about PREPARE, see PREPARE. Also see Section 36.5 for more details about using placeholders and input parameters.

36.4. Using Host Variables

In Section 36.3 you saw how you can execute SQL statements from an embedded SQL program. Some of those statements only used fixed values and did not provide a way to insert user-supplied values into statements or have the program process the values returned by the query. Those kinds of statements are not really useful in real applications. This section explains in detail how you can pass data between your C program and the embedded SQL statements using a simple mechanism called host variables. In an embedded SQL program we consider the SQL statements to be guests in the C program code which is the host language. Therefore the variables of the C program are called host variables.

Another way to exchange values between PostgreSQL backends and ECPG applications is the use of SQL descriptors, described in Section 36.7.

36.4.1. Overview

Passing data between the C program and the SQL statements is particularly simple in embedded SQL. Instead of having the program paste the data into the statement, which entails various complications, such as properly quoting the value, you can simply write the name of a C variable into the SQL statement, prefixed by a colon. For example:

EXEC SQL INSERT INTO sometable VALUES (:v1, 'foo', :v2);

This statement refers to two C variables named v1 and v2 and also uses a regular SQL string literal, to illustrate that you are not restricted to use one kind of data or the other.

This style of inserting C variables in SQL statements works anywhere a value expression is expected in an SQL statement.

36.4.2. Declare Sections

To pass data from the program to the database, for example as parameters in a query, or to pass data from the database back to the program, the C variables that are intended to contain this data need to be declared in specially marked sections, so the embedded SQL preprocessor is made aware of them.

This section starts with:
EXEC SQL BEGIN DECLARE SECTION;

and ends with:

EXEC SQL END DECLARE SECTION;

Between those lines, there must be normal C variable declarations, such as:

```c
int x = 4;
char foo[16], bar[16];
```

As you can see, you can optionally assign an initial value to the variable. The variable’s scope is determined by the location of its declaring section within the program. You can also declare variables with the following syntax which implicitly creates a declare section:

```c
EXEC SQL int i = 4;
```

You can have as many declare sections in a program as you like.

The declarations are also echoed to the output file as normal C variables, so there’s no need to declare them again. Variables that are not intended to be used in SQL commands can be declared normally outside these special sections.

The definition of a structure or union also must be listed inside a DECLARE section. Otherwise the preprocessor cannot handle these types since it does not know the definition.

### 36.4.3. Retrieving Query Results

Now you should be able to pass data generated by your program into an SQL command. But how do you retrieve the results of a query? For that purpose, embedded SQL provides special variants of the usual commands SELECT and FETCH. These commands have a special INTO clause that specifies which host variables the retrieved values are to be stored in. SELECT is used for a query that returns only single row, and FETCH is used for a query that returns multiple rows, using a cursor.

Here is an example:

```sql
/*
 * assume this table:
 * CREATE TABLE test1 (a int, b varchar(50));
 */

EXEC SQL BEGIN DECLARE SECTION;
int v1;
VARCHAR v2;
EXEC SQL END DECLARE SECTION;

EXEC SQL SELECT a, b INTO :v1, :v2 FROM test;

So the INTO clause appears between the select list and the FROM clause. The number of elements in the select list and the list after INTO (also called the target list) must be equal.

Here is an example using the command FETCH:
EXEC SQL BEGIN DECLARE SECTION;
int v1;
VARCHAR v2;
EXEC SQL END DECLARE SECTION;

...

EXEC SQL DECLARE foo CURSOR FOR SELECT a, b FROM test;

...

do
{
  ...
  EXEC SQL FETCH NEXT FROM foo INTO :v1, :v2;
  ...
} while (...);

Here the INTO clause appears after all the normal clauses.

36.4.4. Type Mapping

When ECPG applications exchange values between the PostgreSQL server and the C application, such as when retrieving query results from the server or executing SQL statements with input parameters, the values need to be converted between PostgreSQL data types and host language variable types (C language data types, concretely). One of the main points of ECPG is that it takes care of this automatically in most cases.

In this respect, there are two kinds of data types: Some simple PostgreSQL data types, such as integer and text, can be read and written by the application directly. Other PostgreSQL data types, such as timestamp and numeric can only be accessed through special library functions; see Section 36.4.4.2.

Table 36.1 shows which PostgreSQL data types correspond to which C data types. When you wish to send or receive a value of a given PostgreSQL data type, you should declare a C variable of the corresponding C data type in the declare section.

Table 36.1. Mapping Between PostgreSQL Data Types and C Variable Types

<table>
<thead>
<tr>
<th>PostgreSQL data type</th>
<th>Host variable type</th>
</tr>
</thead>
<tbody>
<tr>
<td>smallint</td>
<td>short</td>
</tr>
<tr>
<td>integer</td>
<td>int</td>
</tr>
<tr>
<td>bigint</td>
<td>long long int</td>
</tr>
<tr>
<td>decimal</td>
<td>decimal(^a)</td>
</tr>
<tr>
<td>numeric</td>
<td>numeric(^a)</td>
</tr>
<tr>
<td>real</td>
<td>float</td>
</tr>
<tr>
<td>double precision</td>
<td>double</td>
</tr>
<tr>
<td>smallserial</td>
<td>short</td>
</tr>
<tr>
<td>serial</td>
<td>int</td>
</tr>
<tr>
<td>bigserial</td>
<td>long long int</td>
</tr>
<tr>
<td>oid</td>
<td>unsigned int</td>
</tr>
<tr>
<td>character(n), varchar(n), text</td>
<td>char[n+1], VARCHAR[n+1](^b)</td>
</tr>
<tr>
<td>name</td>
<td>char[NAMEDATALEN]</td>
</tr>
</tbody>
</table>
### 36.4.4.1. Handling Character Strings

To handle SQL character string data types, such as `varchar` and `text`, there are two possible ways to declare the host variables.

One way is using `char[]`, an array of `char`, which is the most common way to handle character data in C.

```c
EXEC SQL BEGIN DECLARE SECTION;
    char str[50];
EXEC SQL END DECLARE SECTION;
```

Note that you have to take care of the length yourself. If you use this host variable as the target variable of a query which returns a string with more than 49 characters, a buffer overflow occurs.

The other way is using the `VARCHAR` type, which is a special type provided by ECPG. The definition on an array of type `VARCHAR` is converted into a named `struct` for every variable. A declaration like:

```c
VARCHAR var[180];
```

is converted into:

```c
struct varchar_var { int len; char arr[180]; } var;
```

The member `arr` hosts the string including a terminating zero byte. Thus, to store a string in a `VARCHAR` host variable, the host variable has to be declared with the length including the zero byte terminator. The member `len` holds the length of the string stored in the `arr` without the terminating zero byte. When a host variable is used as input for a query, if `strlen(arr)` and `len` are different, the shorter one is used.

`VARCHAR` can be written in upper or lower case, but not in mixed case.

`char` and `VARCHAR` host variables can also hold values of other SQL types, which will be stored in their string forms.

### 36.4.4.2. Accessing Special Data Types

ECPG contains some special types that help you to interact easily with some special data types from the PostgreSQL server. In particular, it has implemented support for the `numeric`, `decimal`, `date`, `timestamp`, and `interval` types. These data types cannot usefully be mapped to primitive host variable types (such as `int`, `long long int`, or `char[]`), because they have a complex internal structure. Applications deal with these types by declaring host variables in special types and accessing them using functions in the `pgtypes` library. The `pgtypes` library, described in detail in Section 36.6 contains basic functions to deal with those types, such that you do not need to send a query to the SQL server just for adding an interval to a time stamp for example.
The follow subsections describe these special data types. For more details about pgtypes library functions, see Section 36.6.

### 36.4.4.2.1. timestamp, date

Here is a pattern for handling timestamp variables in the ECPG host application.

First, the program has to include the header file for the timestamp type:

```
#include <pgtypes_timestamp.h>
```

Next, declare a host variable as type timestamp in the declare section:

```
EXEC SQL BEGIN DECLARE SECTION;  
timestamp ts;  
EXEC SQL END DECLARE SECTION;
```

And after reading a value into the host variable, process it using pgtypes library functions. In following example, the timestamp value is converted into text (ASCII) form with the `PGTYPEStimestamp_to_asc()` function:

```
EXEC SQL SELECT now()::timestamp INTO :ts;  
printf("ts = %s\n", PGTYPEStimestamp_to_asc(ts));
```

This example will show some result like following:

```
ts = 2010-06-27 18:03:56.949343
```

In addition, the DATE type can be handled in the same way. The program has to include pgtypes_date.h, declare a host variable as the date type and convert a DATE value into a text form using `PGTYPESdate_to_asc()` function. For more details about the pgtypes library functions, see Section 36.6.

### 36.4.4.2.2. interval

The handling of the interval type is also similar to the timestamp and date types. It is required, however, to allocate memory for an interval type value explicitly. In other words, the memory space for the variable has to be allocated in the heap memory, not in the stack memory.

Here is an example program:

```
#include <stdio.h>
#include <stdlib.h>
#include <pgtypes_interval.h>

int main(void)  
{  
EXEC SQL BEGIN DECLARE SECTION;  
    interval *in;  
EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO testdb;  
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
```
EXEC SQL SELECT '1 min'::interval INTO :in;
printf("interval = %s\n", PGTYPESinterval_to_asc(in));
PGTYPESinterval_free(in);
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
return 0;
}

36.4.4.2.3. numeric, decimal

The handling of the numeric and decimal types is similar to the interval type: It requires defining a pointer, allocating some memory space on the heap, and accessing the variable using the pgtypes library functions. For more details about the pgtypes library functions, see Section 36.6.

No functions are provided specifically for the decimal type. An application has to convert it to a numeric variable using a pgtypes library function to do further processing.

Here is an example program handling numeric and decimal type variables.

#include <stdio.h>
#include <stdlib.h>
#include <pgtypes_numeric.h>
EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
EXEC SQL BEGIN DECLARE SECTION;
numeric *num;
numeric *num2;
decimal *dec;
EXEC SQL END DECLARE SECTION;
EXEC SQL CONNECT TO testdb;
EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

num = PGTYPESnumeric_new();
dec = PGTYPESdecimal_new();

EXEC SQL SELECT 12.345::numeric(4,2), 23.456::decimal(4,2) INTO :num, :dec;

printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 0));
printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 1));
printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 2));

/* Convert decimal to numeric to show a decimal value. */
num2 = PGTYPESnumeric_new();
PGTYPESnumeric_from_decimal(dec, num2);

printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 0));
printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 1));
printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 2));
ECPG - Embedded SQL in C

36.4.4.3. Host Variables with Nonprimitive Types

As a host variable you can also use arrays, typedefs, structs, and pointers.

36.4.4.3.1. Arrays

There are two use cases for arrays as host variables. The first is a way to store some text string in
char[] or VARCHAR[], as explained in Section 36.4.4.1. The second use case is to retrieve multiple
rows from a query result without using a cursor. Without an array, to process a query result consisting
of multiple rows, it is required to use a cursor and the FETCH command. But with array host variables,
multiple rows can be received at once. The length of the array has to be defined to be able to accom-
modate all rows, otherwise a buffer overflow will likely occur.

Following example scans the pg_database system table and shows all OIDs and names of the
available databases:

```c
int
main(void)
{
  EXEC SQL BEGIN DECLARE SECTION;
  int dbid[8];
  char dbname[8][16];
  int i;
  EXEC SQL END DECLARE SECTION;

  memset(dbname, 0, sizeof(char)* 16 * 8);
  memset(dbid, 0, sizeof(int) * 8);

  EXEC SQL CONNECT TO testdb;
  EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

  /* Retrieve multiple rows into arrays at once. */
  EXEC SQL SELECT oid, datname INTO :dbid, :dbname FROM pg_database;

  for (i = 0; i < 8; i++)
    printf("oid=%d, dbname=%s\n", dbid[i], dbname[i]);

  EXEC SQL COMMIT;
  EXEC SQL DISCONNECT ALL;
  return 0;
}
```

This example shows following result. (The exact values depend on local circumstances.)
36.4.4.3.2. Structures

A structure whose member names match the column names of a query result, can be used to retrieve multiple columns at once. The structure enables handling multiple column values in a single host variable.

The following example retrieves OIDs, names, and sizes of the available databases from the pg_database system table and using the pg_database_size() function. In this example, a structure variable dbinfo_t with members whose names match each column in the SELECT result is used to retrieve one result row without putting multiple host variables in the FETCH statement.

EXEC SQL BEGIN DECLARE SECTION;
  typedef struct
  {
    int oid;
    char datname[65];
    long long int size;
  } dbinfo_t;

  dbinfo_t dbval;
EXEC SQL END DECLARE SECTION;

  memset(&dbval, 0, sizeof(dbinfo_t));

EXEC SQL DECLARE curl CURSOR FOR SELECT oid, datname, pg_database_size(oid) AS size FROM pg_database;
EXEC SQL OPEN curl;

  /* when end of result set reached, break out of while loop */
  EXEC SQL WHENEVER NOT FOUND DO BREAK;

  while (1)
  {
    /* Fetch multiple columns into one structure. */
    EXEC SQL FETCH FROM curl INTO :dbval;

    /* Print members of the structure. */
    printf("oid=%d, datname=%s, size=%lld\n", dbval.oid, dbval.datname, dbval.size);
  }

EXEC SQL CLOSE curl;

This example shows following result. (The exact values depend on local circumstances.)

oid=1, dbname=template1
oid=11510, dbname=template0
oid=11511, dbname=postgres
oid=313780, dbname=testdb

oid=0, dbname=
oid=0, dbname=
oid=0, dbname=

oid=1, datname=template1, size=4324580
oid=11510, datname=template0, size=4243460
oid=11511, datname=postgres, size=4324580
oid=313780, datname=testdb, size=8183012
Structure host variables “absorb” as many columns as the structure as fields. Additional columns can be assigned to other host variables. For example, the above program could also be restructured like this, with the size variable outside the structure:

```c
EXEC SQL BEGIN DECLARE SECTION;
    typedef struct
    {
        int oid;
        char datname[65];
    } dbinfo_t;

dbinfo_t dbval;
long long int size;
EXEC SQL END DECLARE SECTION;

memset(&dbval, 0, sizeof(dbinfo_t));

EXEC SQL DECLARE cur1 CURSOR FOR SELECT oid, datname, pg_database_size(oid) AS size FROM pg_database;
EXEC SQL OPEN cur1;

/* when end of result set reached, break out of while loop */
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
{
    /* Fetch multiple columns into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :dbval, :size;

    /* Print members of the structure. */
    printf("oid=%d, datname=%s, size=%lld\n", dbval.oid, dbval.datname, size);
}
EXEC SQL CLOSE cur1;
```

### 36.4.4.3.3. Typedefs

Use the `typedef` keyword to map new types to already existing types.

```c
EXEC SQL BEGIN DECLARE SECTION;
    typedef char mychartype[40];
    typedef long serial_t;
EXEC SQL END DECLARE SECTION;

Note that you could also use:

```c
EXEC SQL TYPE serial_t IS long;
```

This declaration does not need to be part of a declare section.

### 36.4.4.3.4. Pointers

You can declare pointers to the most common types. Note however that you cannot use pointers as target variables of queries without auto-allocation. See Section 36.7 for more information on auto-allocation.
EXEC SQL BEGIN DECLARE SECTION;
  int *intp;
  char **charp;
EXEC SQL END DECLARE SECTION;

36.4.5. Handling Nonprimitive SQL Data Types

This section contains information on how to handle nonscalar and user-defined SQL-level data types in ECPG applications. Note that this is distinct from the handling of host variables of nonprimitive types, described in the previous section.

36.4.5.1. Arrays

Multi-dimensional SQL-level arrays are not directly supported in ECPG. One-dimensional SQL-level arrays can be mapped into C array host variables and vice-versa. However, when creating a statement ecpg does not know the types of the columns, so that it cannot check if a C array is input into a corresponding SQL-level array. When processing the output of a SQL statement, ecpg has the necessary information and thus checks if both are arrays.

If a query accesses elements of an array separately, then this avoids the use of arrays in ECPG. Then, a host variable with a type that can be mapped to the element type should be used. For example, if a column type is array of integer, a host variable of type int can be used. Also if the element type is varchar or text, a host variable of type char[] or VARCHAR[] can be used.

Here is an example. Assume the following table:

```
CREATE TABLE t3 (  
  ii integer[]  
);  
testdb=> SELECT * FROM t3;
  ii  
-----  
{1,2,3,4,5}  
(1 row)
```

The following example program retrieves the 4th element of the array and stores it into a host variable of type int:

```
EXEC SQL BEGIN DECLARE SECTION;
  int ii;
EXEC SQL END DECLARE SECTION;

EXEC SQL DECLARE curl CURSOR FOR SELECT ii[4] FROM t3;
EXEC SQL OPEN curl;

EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)  
{  
  EXEC SQL FETCH FROM curl INTO :ii ;  
  printf("ii=%d\n", ii);  
}

EXEC SQL CLOSE curl;
```
This example shows the following result:

\[ ii=4 \]

To map multiple array elements to the multiple elements in an array type host variables each element of array column and each element of the host variable array have to be managed separately, for example:

```sql
EXEC SQL BEGIN DECLARE SECTION;
int ii_a[8];
EXEC SQL END DECLARE SECTION;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT ii[1], ii[2], ii[3], ii[4] FROM t3;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    EXEC SQL FETCH FROM cur1 INTO :ii_a[0], :ii_a[1], :ii_a[2], :ii_a[3];
    ...
}
```

Note again that

```sql
EXEC SQL BEGIN DECLARE SECTION;
int ii_a[8];
EXEC SQL END DECLARE SECTION;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT ii FROM t3;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    /* WRONG */
    EXEC SQL FETCH FROM cur1 INTO :ii_a;
    ...
}
```

would not work correctly in this case, because you cannot map an array type column to an array host variable directly.

Another workaround is to store arrays in their external string representation in host variables of type `char[]` or `VARCHAR[]`. For more details about this representation, see Section 8.15.2. Note that this means that the array cannot be accessed naturally as an array in the host program (without further processing that parses the text representation).

### 36.4.5.2. Composite Types

Composite types are not directly supported in ECPG, but an easy workaround is possible. The available workarounds are similar to the ones described for arrays above: Either access each attribute separately or use the external string representation.

For the following examples, assume the following type and table:
CREATE TYPE comp_t AS (intval integer, textval varchar(32));
CREATE TABLE t4 (compval comp_t);
INSERT INTO t4 VALUES ( (256, 'PostgreSQL') );

The most obvious solution is to access each attribute separately. The following program retrieves data from the example table by selecting each attribute of the type comp_t separately:

EXEC SQL BEGIN DECLARE SECTION;
int intval;
varchar textval[33];
EXEC SQL END DECLARE SECTION;

/* Put each element of the composite type column in the SELECT list. */
EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).intval,
   (compval).textval FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
{
   /* Fetch each element of the composite type column into host variables. */
   EXEC SQL FETCH FROM cur1 INTO :intval, :textval;
   printf("intval=%d, textval=%s\n", intval, textval.arr);
}
EXEC SQL CLOSE cur1;

To enhance this example, the host variables to store values in the FETCH command can be gathered into one structure. For more details about the host variable in the structure form, see Section 36.4.4.3.2.

To switch to the structure, the example can be modified as below. The two host variables, intval and textval, become members of the comp_t structure, and the structure is specified on the FETCH command.

EXEC SQL BEGIN DECLARE SECTION;
typedef struct
{
   int intval;
   varchar textval[33];
} comp_t;

comp_t compval;
EXEC SQL END DECLARE SECTION;

/* Put each element of the composite type column in the SELECT list. */
EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).intval,
   (compval).textval FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
{ /* Put all values in the SELECT list into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :compval;

    printf("intval=%d, textval=%s\n", compval.intval,
            compval.textval.arr);
}
EXEC SQL CLOSE cur1;

Although a structure is used in the FETCH command, the attribute names in the SELECT clause are
specified one by one. This can be enhanced by using a * to ask for all attributes of the composite
type value.

... EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).* FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    /* Put all values in the SELECT list into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :compval;

    printf("intval=%d, textval=%s\n", compval.intval,
            compval.textval.arr);
}
...

This way, composite types can be mapped into structures almost seamlessly, even though ECPG does
not understand the composite type itself.

Finally, it is also possible to store composite type values in their external string representation in host
variables of type char[] or VARCHAR[]. But that way, it is not easily possible to access the fields
of the value from the host program.

36.4.5.3. User-defined Base Types

New user-defined base types are not directly supported by ECPG. You can use the external string
representation and host variables of type char[] or VARCHAR[], and this solution is indeed appro-
priate and sufficient for many types.

Here is an example using the data type complex from the example in Section 38.12. The external
string representation of that type is (%f,%f), which is defined in the functions complex_in() and
complex_out() functions in Section 38.12. The following example inserts the complex type
values (1,1) and (3,3) into the columns a and b, and select them from the table after that.

EXEC SQL BEGIN DECLARE SECTION;
    varchar a[64];
    varchar b[64];
EXEC SQL END DECLARE SECTION;

EXEC SQL INSERT INTO test_complex VALUES ('(1,1)', '(3,3)');
EXEC SQL DECLARE cur1 CURSOR FOR SELECT a, b FROM test_complex;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
{
    EXEC SQL FETCH FROM cur1 INTO :a, :b;
    printf("a=%s, b=%s\n", a.arr, b.arr);
}

EXEC SQL CLOSE cur1;

This example shows following result:

a=(1,1), b=(3,3)

Another workaround is avoiding the direct use of the user-defined types in ECPG and instead create a function or cast that converts between the user-defined type and a primitive type that ECPG can handle. Note, however, that type casts, especially implicit ones, should be introduced into the type system very carefully.

For example,

CREATE FUNCTION create_complex(r double, i double) RETURNS complex
LANGUAGE SQL
IMMUTABLE
AS $$ SELECT $1 * complex '(1,0)' + $2 * complex '(0,1)' $$;

After this definition, the following

EXEC SQL BEGIN DECLARE SECTION;
double a, b, c, d;
EXEC SQL END DECLARE SECTION;

a = 1;
b = 2;
c = 3;
d = 4;

EXEC SQL INSERT INTO test_complex VALUES (create_complex(:a, :b),
create_complex(:c, :d));

has the same effect as

EXEC SQL INSERT INTO test_complex VALUES ('(1,2)', '(3,4)');

36.4.6. Indicators

The examples above do not handle null values. In fact, the retrieval examples will raise an error if they fetch a null value from the database. To be able to pass null values to the database or retrieve null values from the database, you need to append a second host variable specification to each host variable that contains data. This second host variable is called the indicator and contains a flag that tells whether the datum is null, in which case the value of the real host variable is ignored. Here is an example that handles the retrieval of null values correctly:
EXEC SQL BEGIN DECLARE SECTION;
  VARCHAR val;
  int val_ind;
EXEC SQL END DECLARE SECTION:

...

EXEC SQL SELECT b INTO :val :val_ind FROM test1;

The indicator variable val_ind will be zero if the value was not null, and it will be negative if the value was null.

The indicator has another function: if the indicator value is positive, it means that the value is not null, but it was truncated when it was stored in the host variable.

If the argument \_r no_indicator is passed to the preprocessor ecpg, it works in “no-indicator” mode. In no-indicator mode, if no indicator variable is specified, null values are signaled (on input and output) for character string types as empty string and for integer types as the lowest possible value for type (for example, INT_MIN for int).

36.5. Dynamic SQL

In many cases, the particular SQL statements that an application has to execute are known at the time the application is written. In some cases, however, the SQL statements are composed at run time or provided by an external source. In these cases you cannot embed the SQL statements directly into the C source code, but there is a facility that allows you to call arbitrary SQL statements that you provide in a string variable.

36.5.1. Executing Statements without a Result Set

The simplest way to execute an arbitrary SQL statement is to use the command EXECUTE IMMEDIATE. For example:

EXEC SQL BEGIN DECLARE SECTION;
  const char *stmt = "CREATE TABLE test1 ( ... );";
EXEC SQL END DECLARE SECTION;

EXEC SQL EXECUTE IMMEDIATE :stmt;

EXECUTE IMMEDIATE can be used for SQL statements that do not return a result set (e.g., DDL, INSERT, UPDATE, DELETE). You cannot execute statements that retrieve data (e.g., SELECT) this way. The next section describes how to do that.

36.5.2. Executing a Statement with Input Parameters

A more powerful way to execute arbitrary SQL statements is to prepare them once and execute the prepared statement as often as you like. It is also possible to prepare a generalized version of a statement and then execute specific versions of it by substituting parameters. When preparing the statement, write question marks where you want to substitute parameters later. For example:

EXEC SQL BEGIN DECLARE SECTION;
  const char *stmt = "INSERT INTO test1 VALUES(?, ?);";
EXEC SQL END DECLARE SECTION;

EXEC SQL PREPARE mystmt FROM :stmt;

...
EXEC SQL EXECUTE mystmt USING 42, 'foobar';

When you don’t need the prepared statement anymore, you should deallocate it:

EXEC SQL DEALLOCATE PREPARE name;

### 36.5.3. Executing a Statement with a Result Set

To execute an SQL statement with a single result row, EXECUTE can be used. To save the result, add an INTO clause.

```sql
EXEC SQL BEGIN DECLARE SECTION;
const char *stmt = "SELECT a, b, c FROM test1 WHERE a > ?";
int v1, v2;
VARCHAR v3[50];
EXEC SQL END DECLARE SECTION;
EXEC SQL PREPARE mystmt FROM :stmt;
...
EXEC SQL EXECUTE mystmt INTO :v1, :v2, :v3 USING 37;
```

An EXECUTE command can have an INTO clause, a USING clause, both, or neither.

If a query is expected to return more than one result row, a cursor should be used, as in the following example. (See Section 36.3.2 for more details about the cursor.)

```sql
EXEC SQL BEGIN DECLARE SECTION;
char dbaname[128];
char datname[128];
char *stmt = "SELECT u.usename as dbaname, d.datname "
"  FROM pg_database d, pg_user u "
"  WHERE d.datdba = u.usesysid";
EXEC SQL END DECLARE SECTION;
EXEC SQL CONNECT TO testdb AS con1 USER testuser;
EXEC SQL SELECT pg_catalog.set_config('search_path', '', false);
EXEC SQL COMMIT;
EXEC SQL PREPARE stmt1 FROM :stmt;
EXEC SQL DECLARE cursor1 CURSOR FOR stmt1;
EXEC SQL OPEN cursor1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1)
{
    EXEC SQL FETCH cursor1 INTO :dbaname,:datname;
    printf("dbaname=%s, datname=%s\n", dbaname, datname);
}
EXEC SQL CLOSE cursor1;
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
```
36.6. pgtypes Library

The pgtypes library maps PostgreSQL database types to C equivalents that can be used in C programs. It also offers functions to do basic calculations with those types within C, i.e., without the help of the PostgreSQL server. See the following example:

```sql
EXEC SQL BEGIN DECLARE SECTION;
   date date1;
   timestamp tsl1, tsout;
   interval iv1;
   char *out;
EXEC SQL END DECLARE SECTION;
PGTYPESdate_today(&date1);
EXEC SQL SELECT started, duration INTO :ts1, :iv1 FROM datetbl
   WHERE d=:date1;
PGTYPEStimestamp_add_interval(&ts1, &iv1, &tsout);
out = PGTYPEStimestamp_to_asc(&tsout);
printf("Started + duration: %s\n", out);
PGTYPESchar_free(out);
```

36.6.1. Character Strings

Some functions such as `PGTYPESnumeric_to_asc` return a pointer to a freshly allocated character string. These results should be freed with `PGTYPESchar_free` instead of `free`. (This is important only on Windows, where memory allocation and release sometimes need to be done by the same library.)

36.6.2. The numeric Type

The numeric type offers to do calculations with arbitrary precision. See Section 8.1 for the equivalent type in the PostgreSQL server. Because of the arbitrary precision this variable needs to be able to expand and shrink dynamically. That's why you can only create numeric variables on the heap, by means of the `PGTYPESnumeric_new` and `PGTYPESnumeric_free` functions. The decimal type, which is similar but limited in precision, can be created on the stack as well as on the heap.

The following functions can be used to work with the numeric type:

```c
PGTYPESnumeric_new
    Request a pointer to a newly allocated numeric variable.

    numeric *PGTYPESnumeric_new(void);

PGTYPESnumeric_free
    Free a numeric type, release all of its memory.

    void PGTYPESnumeric_free(numeric *var);

PGTYPESnumeric_from_asc
    Parse a numeric type from its string notation.
```
numeric *PGTYPESnumeric_from_asc(char *str, char **endptr);

Valid formats are for example: -2, .794, +3.44, 592.49E07 or -32.84e-4. If the value could be parsed successfully, a valid pointer is returned, else the NULL pointer. At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in *endptr. You can safely set endptr to NULL.

PGTYPESnumeric_to_asc

Returns a pointer to a string allocated by malloc that contains the string representation of the numeric type num.

cchar *PGTYPESnumeric_to_asc(numeric *num, int dscale);

The numeric value will be printed with dscale decimal digits, with rounding applied if necessary. The result must be freed with PGTYPESchar_free().

PGTYPESnumeric_add

Add two numeric variables into a third one.

int PGTYPESnumeric_add(numeric *var1, numeric *var2, numeric *result);

The function adds the variables var1 and var2 into the result variable result. The function returns 0 on success and -1 in case of error.

PGTYPESnumeric_sub

Subtract two numeric variables and return the result in a third one.

int PGTYPESnumeric_sub(numeric *var1, numeric *var2, numeric *result);

The function subtracts the variable var2 from the variable var1. The result of the operation is stored in the variable result. The function returns 0 on success and -1 in case of error.

PGTYPESnumeric_mul

Multiply two numeric variables and return the result in a third one.

int PGTYPESnumeric_mul(numeric *var1, numeric *var2, numeric *result);

The function multiplies the variables var1 and var2. The result of the operation is stored in the variable result. The function returns 0 on success and -1 in case of error.

PGTYPESnumeric_div

Divide two numeric variables and return the result in a third one.

int PGTYPESnumeric_div(numeric *var1, numeric *var2, numeric *result);

The function divides the variables var1 by var2. The result of the operation is stored in the variable result. The function returns 0 on success and -1 in case of error.
PGTYPESnumeric_cmp

Compare two numeric variables.

int PGTYPESnumeric_cmp(numeric *var1, numeric *var2)

This function compares two numeric variables. In case of error, INT_MAX is returned. On success, the function returns one of three possible results:

• 1, if var1 is bigger than var2
• -1, if var1 is smaller than var2
• 0, if var1 and var2 are equal

PGTYPESnumeric_from_int

Convert an int variable to a numeric variable.

int PGTYPESnumeric_from_int(signed int int_val, numeric *var);

This function accepts a variable of type signed int and stores it in the numeric variable var. Upon success, 0 is returned and -1 in case of a failure.

PGTYPESnumeric_from_long

Convert a long int variable to a numeric variable.

int PGTYPESnumeric_from_long(signed long int long_val, numeric *var);

This function accepts a variable of type signed long int and stores it in the numeric variable var. Upon success, 0 is returned and -1 in case of a failure.

PGTYPESnumeric_copy

Copy over one numeric variable into another one.

int PGTYPESnumeric_copy(numeric *src, numeric *dst);

This function copies over the value of the variable that src points to into the variable that dst points to. It returns 0 on success and -1 if an error occurs.

PGTYPESnumeric_from_double

Convert a variable of type double to a numeric.

int PGTYPESnumeric_from_double(double d, numeric *dst);

This function accepts a variable of type double and stores the result in the variable that dst points to. It returns 0 on success and -1 if an error occurs.

PGTYPESnumeric_to_double

Convert a variable of type numeric to double.
int PGTYPESnumeric_to_double(numeric *nv, double *dp)

The function converts the numeric value from the variable that \( nv \) points to into the double variable that \( dp \) points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable \( errno \) will be set to PGTYPES_NUM_OVERFLOW additionally.

PGTYPESnumeric_to_int

Convert a variable of type numeric to int.

int PGTYPESnumeric_to_int(numeric *nv, int *ip);

The function converts the numeric value from the variable that \( nv \) points to into the integer variable that \( ip \) points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable \( errno \) will be set to PGTYPES_NUM_OVERFLOW additionally.

PGTYPESnumeric_to_long

Convert a variable of type numeric to long.

int PGTYPESnumeric_to_long(numeric *nv, long *lp);

The function converts the numeric value from the variable that \( nv \) points to into the long integer variable that \( lp \) points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable \( errno \) will be set to PGTYPES_NUM_OVERFLOW additionally.

PGTYPESnumeric_to_decimal

Convert a variable of type numeric to decimal.

int PGTYPESnumeric_to_decimal(numeric *src, decimal *dst);

The function converts the numeric value from the variable that \( src \) points to into the decimal variable that \( dst \) points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable \( errno \) will be set to PGTYPES_NUM_OVERFLOW additionally.

PGTYPESnumeric_from_decimal

Convert a variable of type decimal to numeric.

int PGTYPESnumeric_from_decimal(decimal *src, numeric *dst);

The function converts the decimal value from the variable that \( src \) points to into the numeric variable that \( dst \) points to. It returns 0 on success and -1 if an error occurs. Since the decimal type is implemented as a limited version of the numeric type, overflow cannot occur with this conversion.

### 36.6.3. The date Type

The date type in C enables your programs to deal with data of the SQL type date. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the date type:

PGTYPESdate_from_timestamp

Extract the date part from a timestamp.
date PGTYPESdate_from_timestamp(timestamp dt);

The function receives a timestamp as its only argument and returns the extracted date part from this timestamp.

**PGTYPESdate_from_asc**

Parse a date from its textual representation.

date PGTYPESdate_from_asc(char *str, char **endptr);

The function receives a C char* string `str` and a pointer to a C char* string `endptr`. At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in `endptr`. You can safely set `endptr` to NULL.

Note that the function always assumes MDY-formatted dates and there is currently no variable to change that within ECPG.

Table 36.2 shows the allowed input formats.

**Table 36.2. Valid Input Formats for PGTYPESdate_from_asc**

<table>
<thead>
<tr>
<th>Input</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>January 8, 1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1999-01-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1/8/1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1/18/1999</td>
<td>January 18, 1999</td>
</tr>
<tr>
<td>01/02/03</td>
<td>February 1, 2003</td>
</tr>
<tr>
<td>1999-Jan-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>Jan-08-1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>99-Jan-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-99</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-06</td>
<td>January 8, 2006</td>
</tr>
<tr>
<td>Jan-08-99</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>19990108</td>
<td>ISO 8601; January 8, 1999</td>
</tr>
<tr>
<td>990108</td>
<td>ISO 8601; January 8, 1999</td>
</tr>
<tr>
<td>1999.008</td>
<td>year and day of year</td>
</tr>
<tr>
<td>J2451187</td>
<td>Julian day</td>
</tr>
<tr>
<td>January 8, 99 BC</td>
<td>year 99 before the Common Era</td>
</tr>
</tbody>
</table>

**PGTYPESdate_to_asc**

Return the textual representation of a date variable.

char *PGTYPESdate_to_asc(date dDate);

The function receives the date `dDate` as its only parameter. It will output the date in the form `1999-01-18`, i.e., in the `YYYY-MM-DD` format. The result must be freed with `PGTYPESchar_free()`.
PGTYPESdate_julmdy

Extract the values for the day, the month and the year from a variable of type date.

```c
void PGTYPESdate_julmdy(date d, int *mdy);
```

The function receives the date `d` and a pointer to an array of 3 integer values `mdy`. The variable name indicates the sequential order: `mdy[0]` will be set to contain the number of the month, `mdy[1]` will be set to the value of the day and `mdy[2]` will contain the year.

PGTYPESdate_mdyjul

Create a date value from an array of 3 integers that specify the day, the month and the year of the date.

```c
void PGTYPESdate_mdyjul(int *mdy, date *jdate);
```

The function receives the array of the 3 integers (`mdy`) as its first argument and as its second argument a pointer to a variable of type date that should hold the result of the operation.

PGTYPESdate_dayofweek

Return a number representing the day of the week for a date value.

```c
int PGTYPESdate_dayofweek(date d);
```

The function receives the date variable `d` as its only argument and returns an integer that indicates the day of the week for this date.

• 0 - Sunday
• 1 - Monday
• 2 - Tuesday
• 3 - Wednesday
• 4 - Thursday
• 5 - Friday
• 6 - Saturday

PGTYPESdate_today

Get the current date.

```c
void PGTYPESdate_today(date *d);
```

The function receives a pointer to a date variable (`d`) that it sets to the current date.

PGTYPESdate_fmt_asc

Convert a variable of type date to its textual representation using a format mask.
int PGTYPESdate_fmt_asc(date dDate, char *fmtstring, char *outbuf);

The function receives the date to convert (dDate), the format mask (fmtstring) and the string that will hold the textual representation of the date (outbuf).

On success, 0 is returned and a negative value if an error occurred.

The following literals are the field specifiers you can use:

- **dd** - The number of the day of the month.
- **mm** - The number of the month of the year.
- **yy** - The number of the year as a two digit number.
- **yyyy** - The number of the year as a four digit number.
- **ddd** - The name of the day (abbreviated).
- **mmm** - The name of the month (abbreviated).

All other characters are copied 1:1 to the output string.

Table 36.3 indicates a few possible formats. This will give you an idea of how to use this function. All output lines are based on the same date: November 23, 1959.

<table>
<thead>
<tr>
<th>Format</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>mmdyy</td>
<td>112359</td>
</tr>
<tr>
<td>dddmmyy</td>
<td>231159</td>
</tr>
<tr>
<td>yyyydd</td>
<td>591123</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>59/11/23</td>
</tr>
<tr>
<td>yy mm dd</td>
<td>59 11 23</td>
</tr>
<tr>
<td>yy.mm.dd</td>
<td>59.11.23</td>
</tr>
<tr>
<td>.mm. yyyy.dd.</td>
<td>.11.1959.23.</td>
</tr>
<tr>
<td>mmm. dd, yyyy</td>
<td>Nov. 23, 1959</td>
</tr>
<tr>
<td>mmm dd yyyy</td>
<td>Nov 23 1959</td>
</tr>
<tr>
<td>yyyy dd mm</td>
<td>1959 23 11</td>
</tr>
<tr>
<td>ddd, mmm. dd, yyyy</td>
<td>Mon, Nov. 23, 1959</td>
</tr>
<tr>
<td>(ddd) mmm. dd, yyyy</td>
<td>(Mon) Nov. 23, 1959</td>
</tr>
</tbody>
</table>

PGTYPESdate_defmt_asc

Use a format mask to convert a C char* string to a value of type date.

int PGTYPESdate_defmt_asc(date *d, char *fmt, char *str);

The function receives a pointer to the date value that should hold the result of the operation (d), the format mask to use for parsing the date (fmt) and the C char* string containing the textual representation of the date (str). The textual representation is expected to match the format mask. However you do not need to have a 1:1 mapping of the string to the format mask. The function only analyzes the sequential order and looks for the literals yy or yyyy that indicate the position of the year, mm to indicate the position of the month and dd to indicate the position of the day.

Table 36.4 indicates a few possible formats. This will give you an idea of how to use this function.
### Table 36.4. Valid Input Formats for rdefmtdate

<table>
<thead>
<tr>
<th>Format</th>
<th>String</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>ddmmyy</td>
<td>21-2-54</td>
<td>1954-02-21</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>2-12-54</td>
<td>1954-12-02</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>20111954</td>
<td>1954-11-20</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>130464</td>
<td>1964-04-13</td>
</tr>
<tr>
<td>mmm.dd.yyyy</td>
<td>MAR-12-1967</td>
<td>1967-03-12</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>1954, February 3rd</td>
<td>1954-02-03</td>
</tr>
<tr>
<td>mmm.dd.yyy</td>
<td>041269</td>
<td>1969-04-12</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>In the year 2525, in 2525-07-28 the month of July, mankind will be alive on the 28th day</td>
<td></td>
</tr>
<tr>
<td>dd-mm-yy</td>
<td>I said on the 28th of 2525-07-28 July in the year 2525</td>
<td></td>
</tr>
<tr>
<td>mmm.dd.yyy</td>
<td>9/14/58</td>
<td>1958-09-14</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>47/03/29</td>
<td>1947-03-29</td>
</tr>
<tr>
<td>mmm.dd.yyy</td>
<td>oct 28 1975</td>
<td>1975-10-28</td>
</tr>
<tr>
<td>mmdyy</td>
<td>Nov 14th, 1985</td>
<td>1985-11-14</td>
</tr>
</tbody>
</table>

### 36.6.4. The timestamp Type

The timestamp type in C enables your programs to deal with data of the SQL type timestamp. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the timestamp type:

**PGTYPEStimestamp_from_asc**

Parse a timestamp from its textual representation into a timestamp variable.

```c
timestamp PGTYPEStimestamp_from_asc(char *str, char **endptr);
```

The function receives the string to parse (str) and a pointer to a C char* (endptr). At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in *endptr. You can safely set endptr to NULL.

The function returns the parsed timestamp on success. On error, PGTYPESInvalidTimestamp is returned and errno is set to PGTYPES_TS_BAD_TIMESTAMP. See PGTYPESInvalidTimestamp for important notes on this value.

In general, the input string can contain any combination of an allowed date specification, a whitespace character and an allowed time specification. Note that time zones are not supported by ECPG. It can parse them but does not apply any calculation as the PostgreSQL server does for example. Timezone specifiers are silently discarded.

Table 36.5 contains a few examples for input strings.

### Table 36.5. Valid Input Formats for PGTYPEStimestamp_from_asc

<table>
<thead>
<tr>
<th>Input</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>1999-01-08 04:05:06</td>
<td>1999-01-08 04:05:06</td>
</tr>
</tbody>
</table>
### Input

<table>
<thead>
<tr>
<th>Input</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>January 8 04:05:06 1999 PST</td>
<td>1999-01-08 04:05:06</td>
</tr>
<tr>
<td>1999-Jan-08 04:05:06.789-8</td>
<td>1999-01-08 04:05:06.789 (time zone specifier ignored)</td>
</tr>
<tr>
<td>J2451187 04:05-08:00</td>
<td>1999-01-08 04:05:00 (time zone specifier ignored)</td>
</tr>
</tbody>
</table>

### PGTYPEStimestamp_to_asc

Converts a date to a C char* string.

```c
char *PGTYPEStimestamp_to_asc(timestamp tstamp);
```

The function receives the timestamp `tstamp` as its only argument and returns an allocated string that contains the textual representation of the timestamp. The result must be freed with `PGTYPESchar_free()`.

### PGTYPEStimestamp_current

Retrieve the current timestamp.

```c
void PGTYPEStimestamp_current(timestamp *ts);
```

The function retrieves the current timestamp and saves it into the timestamp variable that `ts` points to.

### PGTYPEStimestamp_fmt_asc

Convert a timestamp variable to a C char* using a format mask.

```c
int PGTYPEStimestamp_fmt_asc(timestamp *ts, char *output, int str_len, char *fmtstr);
```

The function receives a pointer to the timestamp to convert as its first argument (`ts`), a pointer to the output buffer (`output`), the maximal length that has been allocated for the output buffer (`str_len`) and the format mask to use for the conversion (`fmtstr`).

Upon success, the function returns 0 and a negative value if an error occurred.

You can use the following format specifiers for the format mask. The format specifiers are the same ones that are used in the `strftime` function in libc. Any non-format specifier will be copied into the output buffer.

- `%A` - is replaced by national representation of the full weekday name.
- `%a` - is replaced by national representation of the abbreviated weekday name.
- `%B` - is replaced by national representation of the full month name.
- `%b` - is replaced by national representation of the abbreviated month name.
- `%C` - is replaced by `(year / 100)` as decimal number; single digits are preceded by a zero.
- `%c` - is replaced by national representation of time and date.
- `%D` - is equivalent to `%m/%d/%y`. 

---

892
• %d - is replaced by the day of the month as a decimal number (01-31).

• %E* %O* - POSIX locale extensions. The sequences %Ec %Ee %Ex %Ey %Od %Oe %OH %Oi %Om %OM %OS %Ou %OU %OV %Ow %OW %Oy are supposed to provide alternative representations.

Additionally %OB implemented to represent alternative months names (used standalone, without day mentioned).

• %e - is replaced by the day of month as a decimal number (1-31); single digits are preceded by a blank.

• %F - is equivalent to %Y-%m-%d.

• %G - is replaced by a year as a decimal number with century. This year is the one that contains the greater part of the week (Monday as the first day of the week).

• %g - is replaced by the same year as in %G, but as a decimal number without century (00-99).

• %H - is replaced by the hour (24-hour clock) as a decimal number (00-23).

• %h - the same as %b.

• %I - is replaced by the hour (12-hour clock) as a decimal number (01-12).

• %j - is replaced by the day of the year as a decimal number (001-366).

• %k - is replaced by the hour (24-hour clock) as a decimal number (0-23); single digits are preceded by a blank.

• %l - is replaced by the hour (12-hour clock) as a decimal number (1-12); single digits are preceded by a blank.

• %M - is replaced by the minute as a decimal number (00-59).

• %m - is replaced by the month as a decimal number (01-12).

• %n - is replaced by a newline.

• %O* - the same as %E*.

• %p - is replaced by national representation of either “ante meridiem” or “post meridiem” as appropriate.

• %R - is equivalent to %H: %M.

• %r - is equivalent to %I: %M: %S %p.

• %S - is replaced by the second as a decimal number (00-60).

• %s - is replaced by the number of seconds since the Epoch, UTC.

• %T - is equivalent to %H: %M: %S

• %t - is replaced by a tab.

• %U - is replaced by the week number of the year (Sunday as the first day of the week) as a decimal number (00-53).

• %u - is replaced by the weekday (Monday as the first day of the week) as a decimal number (1-7).
• %V - is replaced by the week number of the year (Monday as the first day of the week) as a decimal number (01-53). If the week containing January 1 has four or more days in the new year, then it is week 1; otherwise it is the last week of the previous year, and the next week is week 1.

• %v - is equivalent to %e-%b-%Y.

• %W - is replaced by the week number of the year (Monday as the first day of the week) as a decimal number (00-53).

• %w - is replaced by the weekday (Sunday as the first day of the week) as a decimal number (0-6).

• %X - is replaced by national representation of the time.

• %x - is replaced by national representation of the date.

• %Y - is replaced by the year with century as a decimal number.

• %y - is replaced by the year without century as a decimal number (00-99).

• %Z - is replaced by the time zone name.

• %z - is replaced by the time zone offset from UTC; a leading plus sign stands for east of UTC, a minus sign for west of UTC, hours and minutes follow with two digits each and no delimiter between them (common form for RFC 822 date headers).

• %+ - is replaced by national representation of the date and time.

• %-* - GNU libc extension. Do not do any padding when performing numerical outputs.

• $*_ - GNU libc extension. Explicitly specify space for padding.

• %0* - GNU libc extension. Explicitly specify zero for padding.

• % - is replaced by %.

PGTYPEStimestamp_sub

Subtract one timestamp from another one and save the result in a variable of type interval.

```c
int PGTYPEStimestamp_sub(timestamp *ts1, timestamp *ts2, interval *iv);
```

The function will subtract the timestamp variable that ts2 points to from the timestamp variable that ts1 points to and will store the result in the interval variable that iv points to.

Upon success, the function returns 0 and a negative value if an error occurred.

PGTYPEStimestamp_defmt_asc

Parse a timestamp value from its textual representation using a formatting mask.

```c
int PGTYPEStimestamp_defmt_asc(char *str, char *fmt, timestamp *d);
```

The function receives the textual representation of a timestamp in the variable str as well as the formatting mask to use in the variable fmt. The result will be stored in the variable that d points to.
If the formatting mask \( \text{fmt} \) is NULL, the function will fall back to the default formatting mask which is %Y-%m-%d %H:%M:%S.

This is the reverse function to \texttt{PGTYPEStimestamp\_fmt\_asc}. See the documentation there in order to find out about the possible formatting mask entries.

\texttt{PGTYPEStimestamp\_add\_interval}

Add an interval variable to a timestamp variable.

\begin{verbatim}
int PGTYPEStimestamp_add_interval(timestamp *tin, interval *span, timestamp *tout);
\end{verbatim}

The function receives a pointer to a timestamp variable \( \text{tin} \) and a pointer to an interval variable \( \text{span} \). It adds the interval to the timestamp and saves the resulting timestamp in the variable that \( \text{tout} \) points to.

Upon success, the function returns 0 and a negative value if an error occurred.

\texttt{PGTYPEStimestamp\_sub\_interval}

Subtract an interval variable from a timestamp variable.

\begin{verbatim}
int PGTYPEStimestamp_sub_interval(timestamp *tin, interval *span, timestamp *tout);
\end{verbatim}

The function subtracts the interval variable that \( \text{span} \) points to from the timestamp variable that \( \text{tin} \) points to and saves the result into the variable that \( \text{tout} \) points to.

Upon success, the function returns 0 and a negative value if an error occurred.

### 36.6.5. The interval Type

The interval type in C enables your programs to deal with data of the SQL type interval. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the interval type:

\texttt{PGTYPESinterval\_new}

Return a pointer to a newly allocated interval variable.

\begin{verbatim}
interval *PGTYPESinterval_new(void);
\end{verbatim}

\texttt{PGTYPESinterval\_free}

Release the memory of a previously allocated interval variable.

\begin{verbatim}
void PGTYPESinterval_new(interval *intvl);
\end{verbatim}

\texttt{PGTYPESinterval\_from\_asc}

Parse an interval from its textual representation.
interval *PGTYPESinterval_from_asc(char *str, char **endptr);

The function parses the input string `str` and returns a pointer to an allocated interval variable. At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in `*endptr`. You can safely set `endptr` to NULL.

**PGTYPESinterval_to_asc**

Convert a variable of type interval to its textual representation.

char *PGTYPESinterval_to_asc(interval *span);

The function converts the interval variable that `span` points to into a C char*. The output looks like this example: 0 1 day 12 hours 59 mins 10 secs. The result must be freed with `PGTYPESchar_free()`.

**PGTYPESinterval_copy**

Copy a variable of type interval.

int PGTYPESinterval_copy(interval *intvlsrc, interval *intvldest);

The function copies the interval variable that `intvlsrc` points to into the variable that `intvldest` points to. Note that you need to allocate the memory for the destination variable before.

### 36.6.6. The decimal Type

The decimal type is similar to the numeric type. However it is limited to a maximum precision of 30 significant digits. In contrast to the numeric type which can be created on the heap only, the decimal type can be created either on the stack or on the heap (by means of the functions `PGTYPESdecimal_new` and `PGTYPESdecimal_free`). There are a lot of other functions that deal with the decimal type in the Informix compatibility mode described in Section 36.15.

The following functions can be used to work with the decimal type and are not only contained in the `libcompat` library.

**PGTYPESdecimal_new**

Request a pointer to a newly allocated decimal variable.

decimal *PGTYPESdecimal_new(void);

**PGTYPESdecimal_free**

Free a decimal type, release all of its memory.

void PGTYPESdecimal_free(decimal *var);

### 36.6.7. errno Values of pgtypeslib

**PGTYPES_NUM_BAD_NUMERIC**

An argument should contain a numeric variable (or point to a numeric variable) but in fact its in-memory representation was invalid.
An overflow occurred. Since the numeric type can deal with almost arbitrary precision, converting a numeric variable into other types might cause overflow.

An underflow occurred. Since the numeric type can deal with almost arbitrary precision, converting a numeric variable into other types might cause underflow.

A division by zero has been attempted.

An invalid date string was passed to the `PGTYPESdate_from_asc` function.

Invalid arguments were passed to the `PGTYPESdate_defmt_asc` function.

An invalid token in the input string was found by the `PGTYPESdate_defmt_asc` function.

An invalid interval string was passed to the `PGTYPESinterval_from_asc` function, or an invalid interval value was passed to the `PGTYPESinterval_to_asc` function.

There was a mismatch in the day/month/year assignment in the `PGTYPESdate_defmt_asc` function.

An invalid day of the month value was found by the `PGTYPESdate_defmt_asc` function.

An invalid month value was found by the `PGTYPESdate_defmt_asc` function.

An invalid timestamp string was passed to the `PGTYPEStimestamp_from_asc` function, or an invalid timestamp value was passed to the `PGTYPEStimestamp_to_asc` function.

An infinite timestamp value was encountered in a context that cannot handle it.

36.6.8. Special Constants of pgtypeslib

`PGTYPESInvalidTimestamp`

A value of type timestamp representing an invalid time stamp. This is returned by the function `PGTYPEStimestamp_from_asc` on parse error. Note that due to the internal representation...
of the timestamp data type, PGTYPESInvalidTimestamp is also a valid timestamp at the same time. It is set to 1899-12-31 23:59:59. In order to detect errors, make sure that your application does not only test for PGTYPESInvalidTimestamp but also for errno != 0 after each call to PGTYPEStimestamp_from_asc.

36.7. Using Descriptor Areas

An SQL descriptor area is a more sophisticated method for processing the result of a SELECT, FETCH or a DESCRIBE statement. An SQL descriptor area groups the data of one row of data together with metadata items into one data structure. The metadata is particularly useful when executing dynamic SQL statements, where the nature of the result columns might not be known ahead of time. PostgreSQL provides two ways to use Descriptor Areas: the named SQL Descriptor Areas and the C-structure SQLDAs.

36.7.1. Named SQL Descriptor Areas

A named SQL descriptor area consists of a header, which contains information concerning the entire descriptor, and one or more item descriptor areas, which basically each describe one column in the result row.

Before you can use an SQL descriptor area, you need to allocate one:

EXEC SQL ALLOCATE DESCRIPTOR identifier;

The identifier serves as the “variable name” of the descriptor area. When you don’t need the descriptor anymore, you should deallocate it:

EXEC SQL DEALLOCATE DESCRIPTOR identifier;

To use a descriptor area, specify it as the storage target in an INTO clause, instead of listing host variables:

EXEC SQL FETCH NEXT FROM mycursor INTO SQL DESCRIPTOR mydesc;

If the result set is empty, the Descriptor Area will still contain the metadata from the query, i.e. the field names.

For not yet executed prepared queries, the DESCRIBE statement can be used to get the metadata of the result set:

EXEC SQL BEGIN DECLARE SECTION;
char *sql_stmt = "SELECT * FROM table1";
EXEC SQL END DECLARE SECTION;

EXEC SQL PREPARE stmt1 FROM :sql_stmt;
EXEC SQL DESCRIBE stmt1 INTO SQL DESCRIPTOR mydesc;

Before PostgreSQL 9.0, the SQL keyword was optional, so using DESCRIPTOR and SQL DESCRIPTOR produced named SQL Descriptor Areas. Now it is mandatory, omitting the SQL keyword produces SQLDA Descriptor Areas, see Section 36.7.2.

In DESCRIBE and FETCH statements, the INTO and USING keywords can be used to similarly: they produce the result set and the metadata in a Descriptor Area.
Now how do you get the data out of the descriptor area? You can think of the descriptor area as a structure with named fields. To retrieve the value of a field from the header and store it into a host variable, use the following command:

```sql
EXEC SQL GET DESCRIPTOR name :hostvar = field;
```

Currently, there is only one header field defined: `COUNT`, which tells how many item descriptor areas exist (that is, how many columns are contained in the result). The host variable needs to be of an integer type. To get a field from the item descriptor area, use the following command:

```sql
EXEC SQL GET DESCRIPTOR name VALUE num :hostvar = field;
```

`num` can be a literal integer or a host variable containing an integer. Possible fields are:

- **CARDINALITY** (integer)
  - number of rows in the result set

- **DATA**
  - actual data item (therefore, the data type of this field depends on the query)

- **DATETIME_INTERVAL_CODE** (integer)
  - When `TYPE` is 9, `DATETIME_INTERVAL_CODE` will have a value of 1 for `DATE`, 2 for `TIME`, 3 for `TIMESTAMP`, 4 for `TIME WITH TIME ZONE`, or 5 for `TIMESTAMP WITH TIME ZONE`.

- **DATETIME_INTERVAL_PRECISION** (integer)
  - not implemented

- **INDICATOR** (integer)
  - the indicator (indicating a null value or a value truncation)

- **KEY_MEMBER** (integer)
  - not implemented

- **LENGTH** (integer)
  - length of the datum in characters

- **NAME** (string)
  - name of the column

- **NULLABLE** (integer)
  - not implemented

- **OCTET_LENGTH** (integer)
  - length of the character representation of the datum in bytes

- **PRECISION** (integer)
  - precision (for type numeric)
EXECUTE, DECLARE and OPEN statements, the effect of the INTO and USING keywords are different. A Descriptor Area can also be manually built to provide the input parameters for a query or a cursor and using SQL DESCRIPTOR name is the way to pass the input parameters into a parameterized query. The statement to build a named SQL Descriptor Area is below:

```
EXEC SQL SET DESCRIPTOR name VALUE num field = :hostvar;
```

PostgreSQL supports retrieving more that one record in one FETCH statement and storing the data in host variables in this case assumes that the variable is an array. E.g.:

```
EXEC SQL BEGIN DECLARE SECTION;
  int id[5];
EXEC SQL END DECLARE SECTION;
EXEC SQL FETCH 5 FROM mycursor INTO DESCRIPTOR mysqlda;
EXEC SQL GET DESCRIPTOR mysqlda VALUE 1 :id = DATA;
```

### 36.7.2. SQLDA Descriptor Areas

An SQLDA Descriptor Area is a C language structure which can be also used to get the result set and the metadata of a query. One structure stores one record from the result set.

```
EXEC SQL include sqlda.h;
sqlnda_t *mysqlda;
EXEC SQL FETCH 3 FROM mycursor INTO DESCRIPTOR mysqlda;
```

Note that the SQL keyword is omitted. The paragraphs about the use cases of the INTO and USING keywords in Section 36.7.1 also apply here with an addition. In a DESCRIBE statement the DESCRIPTOR keyword can be completely omitted if the INTO keyword is used:

```
EXEC SQL DESCRIBE prepared_statement INTO mysqlda;
```

The general flow of a program that uses SQDLDA is:

1. Prepare a query, and declare a cursor for it.
2. Declare an SQLDA for the result rows.
3. Declare an SQLDA for the input parameters, and initialize them (memory allocation, parameter settings).
4. Open a cursor with the input SQLDA.
5. Fetch rows from the cursor, and store them into an output SQLDA.
6. Read values from the output SQLDA into the host variables (with conversion if necessary).
7. Close the cursor.
8. Free the memory area allocated for the input SQLDA.

36.7.2.1. SQLDA Data Structure

SQLDA uses three data structure types: sqlda_t, sqlvar_t, and struct sqlname.

**Tip**

PostgreSQL's SQLDA has a similar data structure to the one in IBM DB2 Universal Database, so some technical information on DB2's SQLDA could help understanding PostgreSQL's one better.

36.7.2.1.1. sqlda_t Structure

The structure type sqlda_t is the type of the actual SQLDA. It holds one record. And two or more sqlda_t structures can be connected in a linked list with the pointer in the desc_next field, thus representing an ordered collection of rows. So, when two or more rows are fetched, the application can read them by following the desc_next pointer in each sqlda_t node.

The definition of sqlda_t is:

```c
struct sqlda_struct {
    char            sqldaid[8];
    long            sqldabc;
    short           sqln;
    short           sqld;
    struct sqlda_struct *desc_next;
    struct sqlvar_struct sqlvar[1];
};
```

typedef struct sqlda_struct sqlda_t;

The meaning of the fields is:

- **sqldaid**
  - It contains the literal string "SQLDA ".

- **sqldabc**
  - It contains the size of the allocated space in bytes.

- **sqln**
  - It contains the number of input parameters for a parameterized query in case it's passed into OPEN, DECLARE or EXECUTE statements using the USING keyword. In case it's used as output of SELECT, EXECUTE or FETCH statements, its value is the same as sqld statement
sqld

It contains the number of fields in a result set.

desc_next

If the query returns more than one record, multiple linked SQLDA structures are returned, and
desc_next holds a pointer to the next entry in the list.

sqlvar

This is the array of the columns in the result set.

### 36.7.2.1.2. sqlvar_t Structure

The structure type sqlvar_t holds a column value and metadata such as type and length. The definition of the type is:

```c
struct sqlvar_struct {
    short sqltype;
    short sqllen;
    char *sqldata;
    short *sqlind;
    struct sqlname sqlname;
};
typedef struct sqlvar_struct sqlvar_t;
```

The meaning of the fields is:

**sqltype**

Contains the type identifier of the field. For values, see `enum ECPGttype` in `ecpgtype.h`.

**sqllen**

Contains the binary length of the field. e.g. 4 bytes for `ECPGt_int`.

**sqldata**

Points to the data. The format of the data is described in Section 36.4.4.

**sqlind**

Points to the null indicator. 0 means not null, -1 means null.

**sqlname**

The name of the field.

### 36.7.2.1.3. struct sqlName Structure

A struct sqlname structure holds a column name. It is used as a member of the sqlvar_t structure. The definition of the structure is:

```c
#define NAMEDATALEN 64
```
struct sqlname
{
    short length;
    char data[NAMEDATALEN];
};

The meaning of the fields is:

length
    Contains the length of the field name.

data
    Contains the actual field name.

36.7.2.2. Retrieving a Result Set Using an SQLDA

The general steps to retrieve a query result set through an SQLDA are:

1. Declare an sqlda_t structure to receive the result set.
2. Execute FETCH/EXECUTE/DESCRIBE commands to process a query specifying the declared SQLDA.
3. Check the number of records in the result set by looking at sqln, a member of the sqlda_t structure.
4. Get the values of each column from sqlvar[0], sqlvar[1], etc., members of the sqlda_t structure.
5. Go to next row (sqlda_t structure) by following the desc_next pointer, a member of the sqlda_t structure.
6. Repeat above as you need.

Here is an example retrieving a result set through an SQLDA.

First, declare a sqlda_t structure to receive the result set.

sqlda_t *sqlda1;

Next, specify the SQLDA in a command. This is a FETCH command example.

EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;

Run a loop following the linked list to retrieve the rows.

sqlda_t *cur_sqlda;

for (cur_sqlda = sqlda1;
    cur_sqlda != NULL;
    cur_sqlda = cur_sqlda->desc_next)
{
    ...
Inside the loop, run another loop to retrieve each column data (sqlvar_t structure) of the row.

```c
for (i = 0; i < cur_sqlda->sqld; i++)
{
    sqlvar_t v = cur_sqlda->sqlvar[i];
    char *sqldata = v.sqldata;
    short sqllen  = v.sqllen;
    ...
}
```

To get a column value, check the sqltype value, a member of the sqlvar_t structure. Then, switch to an appropriate way, depending on the column type, to copy data from the sqlvar field to a host variable.

```c
char var_buf[1024];
switch (v.sqltype)
{
    case ECPGt_char:
        memset(&var_buf, 0, sizeof(var_buf));
        memcpy(&var_buf, sqldata, (sizeof(var_buf) <= sqllen ? sizeof(var_buf) - 1 : sqllen));
        break;

    case ECPGt_int: /* integer */
        memcpy(&intval, sqldata, sqllen);
        snprintf(var_buf, sizeof(var_buf), "%d", intval);
        break;

    ...
}
```

### 36.7.2.3. Passing Query Parameters Using an SQLDA

The general steps to use an SQLDA to pass input parameters to a prepared query are:

1. Create a prepared query (prepared statement)
2. Declare a sqlda_t structure as an input SQLDA.
3. Allocate memory area (as sqlda_t structure) for the input SQLDA.
4. Set (copy) input values in the allocated memory.
5. Open a cursor with specifying the input SQLDA.

Here is an example.

First, create a prepared statement.

```sql
EXEC SQL BEGIN DECLARE SECTION;
char query[1024] = "SELECT d.oid, * FROM pg_database d,
    pg_stat_database s WHERE d.oid = s.datid AND (d.datname = ? OR
d.oid = ?)");
EXEC SQL END DECLARE SECTION;
```
EXEC SQL PREPARE stmt1 FROM :query;

Next, allocate memory for an SQLDA, and set the number of input parameters in sqln, a member variable of the sqlda_t structure. When two or more input parameters are required for the prepared query, the application has to allocate additional memory space which is calculated by (nr. of params - 1) * sizeof(sqlvar_t). The example shown here allocates memory space for two input parameters.

sqlda_t *sqlda2;

sqlda2 = (sqlda_t *) malloc(sizeof(sqlda_t) + sizeof(sqlvar_t));
memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));

sqlda2->sqln = 2; /* number of input variables */

After memory allocation, store the parameter values into the sqlvar[] array. (This is same array used for retrieving column values when the SQLDA is receiving a result set.) In this example, the input parameters are "postgres", having a string type, and 1, having an integer type.

sqlda2->sqlvar[0].sqltype = ECPGt_char;
sqlda2->sqlvar[0].sqldata = "postgres";
sqlda2->sqlvar[0].sqllen  = 8;

int intval = 1;
sqlda2->sqlvar[1].sqltype = ECPGt_int;
sqlda2->sqlvar[1].sqldata = (char *) &intval;
sqlda2->sqlvar[1].sqllen  = sizeof(intval);

By opening a cursor and specifying the SQLDA that was set up beforehand, the input parameters are passed to the prepared statement.

EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;

Finally, after using input SQLDAs, the allocated memory space must be freed explicitly, unlike SQLDAs used for receiving query results.

free(sqlda2);

36.7.2.4. A Sample Application Using SQLDA

Here is an example program, which describes how to fetch access statistics of the databases, specified by the input parameters, from the system catalogs.

This application joins two system tables, pg_database and pg_stat_database on the database OID, and also fetches and shows the database statistics which are retrieved by two input parameters (a database postgres, and OID 1).

First, declare an SQLDA for input and an SQLDA for output.

EXEC SQL include sqlda.h;

sqlda_t *sqlda1; /* an output descriptor */
sqlda_t *sqlda2; /* an input descriptor */

Next, connect to the database, prepare a statement, and declare a cursor for the prepared statement.
int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    char query[1024] = "SELECT d.oid,* FROM pg_database d,
    pg_stat_database s WHERE d.oid=s.datid AND ( d.datname=? OR
d.oid=? )";
    EXEC SQL END DECLARE SECTION;
    EXEC SQL CONNECT TO testdb AS con1 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
    EXEC SQL PREPARE stmt1 FROM :query;
    EXEC SQL DECLARE cur1 CURSOR FOR stmt1;

    Next, put some values in the input SQLDA for the input parameters. Allocate memory for the input
    SQLDA, and set the number of input parameters to sqln. Store type, value, and value length into
    sqltype, sqldata, and sqllen in the sqlvar structure.

    /* Create SQLDA structure for input parameters. */
    sqlda2 = (sqlda_t *) malloc(sizeof(sqlda_t) +
    sizeof(sqlvar_t));
    memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));
    sqlda2->sqln = 2; /* number of input variables */
    sqlda2->sqlvar[0].sqltype = ECPGt_char;
    sqlda2->sqlvar[0].sqldata = "postgres";
    sqlda2->sqlvar[0].sqllen  = 8;

    intval = 1;
    sqlda2->sqlvar[1].sqltype = ECPGt_int;
    sqlda2->sqlvar[1].sqldata = (char *)&intval;
    sqlda2->sqlvar[1].sqllen  = sizeof(intval);

    After setting up the input SQLDA, open a cursor with the input SQLDA.

    /* Open a cursor with input parameters. */
    EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;

    Fetch rows into the output SQLDA from the opened cursor. (Generally, you have to call FETCH
    repeatedly in the loop, to fetch all rows in the result set.)

    while (1)
    {
        sqlda_t *cur_sqlda;

        /* Assign descriptor to the cursor */
        EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;

        Next, retrieve the fetched records from the SQLDA, by following the linked list of the sqlda_t
        structure.

        for (cur_sqlda = sqlda1 ;
cur_sqlda != NULL ;
cur_sqlda = cur_sqlda->desc_next)
{
  ...
}

Read each columns in the first record. The number of columns is stored in sqld, the actual data of the first column is stored in sqlvar[0], both members of the sqlda_t structure.

/* Print every column in a row. */
for (i = 0; i < sqlda1->sqld; i++)
{
    sqlvar_t v = sqlda1->sqlvar[i];
    char *sqldata = v.sqldata;
    short sqllen = v.sqllen;

    strncpy(name_buf, v.sqlname.data, v.sqlname.length);
    name_buf[v.sqlname.length] = '\0';

Now, the column data is stored in the variable v. Copy every datum into host variables, looking at v.sqltype for the type of the column.

switch (v.sqltype) {
  int intval;
  double doubleval;
  unsigned long long int longlongval;

  case ECPGt_char:
    memset(&var_buf, 0, sizeof(var_buf));
    memcpy(&var_buf, sqldata, (sizeof(var_buf) <= sqllen ? sizeof(var_buf)-1 : sqllen));
    break;

  case ECPGt_int: /* integer */
    memcpy(&intval, sqldata, sqllen);
    snprintf(var_buf, sizeof(var_buf), "%d", intval);
    break;

  ...

  default:
    ...
}

printf("%s = %s (type: %d)\n", name_buf, var_buf, v.sqltype);
}

Close the cursor after processing all of records, and disconnect from the database.

EXEC SQL CLOSE cur1;
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;

The whole program is shown in Example 36.1.
Example 36.1. Example SQLDA Program

```c
#include <stdlib.h>
#include <string.h>
#include <stdlib.h>
#include <stdio.h>
#include <unistd.h>
EXEC SQL include sqlda.h;

sqlda_t *sqlda1; /* descriptor for output */
sqlda_t *sqlda2; /* descriptor for input */
EXEC SQL WHENEVER NOT FOUND DO BREAK;
EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    char query[1024] = "SELECT d.oid,* FROM pg_database d,
pg_stat_database s WHERE d.oid=s.datid AND ( d.datname=? OR
d.oid=? )";

    int intval;
    unsigned long long int longlongval;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO uptimedb AS con1 USER uptime;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '',
false); EXEC SQL COMMIT;

    EXEC SQL PREPARE stmt1 FROM :query;
    EXEC SQL DECLARE cur1 CURSOR FOR stmt1;

    /* Create a SQLDA structure for an input parameter */
    sqlda2 = (sqlda_t *)malloc(sizeof(sqlda_t) + sizeof(sqlvar_t));
    memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));
    sqlda2->sqln = 2; /* a number of input variables */

    sqlda2->sqlvar[0].sqltype = ECPGt_char;
    sqlda2->sqlvar[0].sqldata = "postgres";
    sqlda2->sqlvar[0].sqllen = 8;

    intval = 1;
    sqlda2->sqlvar[1].sqltype = ECPGt_int;
    sqlda2->sqlvar[1].sqldata = (char *) &intval;
    sqlda2->sqlvar[1].sqllen = sizeof(intval);

    /* Open a cursor with input parameters. */
    EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;

    while (1)
    {
        sqlda_t *cur_sqlda;

        /* Assign descriptor to the cursor */
```
EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;

for (cur_sqlda = sqlda1 ;
    cur_sqlda != NULL ;
    cur_sqlda = cur_sqlda->desc_next)
{
    int i;
    char name_buf[1024];
    char var_buf[1024];

    /* Print every column in a row. */
    for (i=0 ; i<cur_sqlda->sqli ; i++)
    {
        sqlvar_t v = cur_sqlda->sqlvar[i];
        char *sqldata = v.sqldata;
        short sqllen = v.sqllen;

        strncpy(name_buf, v.sqlname.data,
                v.sqlname.length);
        name_buf[v.sqlname.length] = '\0';

        switch (v.sqltype)
        {
        case ECPGt_char:
            memset(&var_buf, 0, sizeof(var_buf));
            memcpy(&var_buf, sqldata,
                    (sizeof(var_buf)<=sqllen ? sizeof(var_buf)-1 : sqllen) );
            break;

        case ECPGt_int: /* integer */
            memcpy(&intval, sqldata, sqllen);
            snprintf(var_buf, sizeof(var_buf), "%d",
                    intval);
            break;

        case ECPGt_long_long: /* bigint */
            memcpy(&longlongval, sqldata, sqllen);
            snprintf(var_buf, sizeof(var_buf), "%lld",
                    longlongval);
            break;

        default:
        {
            int i;
            memset(var_buf, 0, sizeof(var_buf));
            for (i = 0; i < sqllen; i++)
            {
                char tmpbuf[16];
                snprintf(tmpbuf, sizeof(tmpbuf), "%02x",
                        (unsigned char) sqldata[i]);
                strncat(var_buf, tmpbuf,
                        sizeof(var_buf));
            }
            break;
        }
    }
}

909
EXEC SQL CLOSE cur1;
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
return 0;
}

The output of this example should look something like the following (some numbers will vary).

oid = 1 (type: 1)
datname = template1 (type: 1)
datdba = 10 (type: 1)
encoding = 0 (type: 5)
datistemplate = t (type: 1)
datallowconn = t (type: 1)
datconnlimit = -1 (type: 5)
datlastsysoid = 11510 (type: 1)
datfrozenxid = 379 (type: 1)
dattablespace = 1663 (type: 1)
datconfig = (type: 1)
datacl = {=c/uptime, uptime=CTc/uptime} (type: 1)
datid = 1 (type: 1)
datname = template1 (type: 1)
numbackends = 0 (type: 5)
xact_commit = 113606 (type: 9)
xact_rollback = 0 (type: 9)
blks_read = 130 (type: 9)
blks_hit = 7341714 (type: 9)
tup_returned = 38262679 (type: 9)
tup_fetched = 1836281 (type: 9)
tup_inserted = 0 (type: 9)
tup_updated = 0 (type: 9)
tup_deleted = 0 (type: 9)

oid = 11511 (type: 1)
datname = postgres (type: 1)
datdba = 10 (type: 1)
encoding = 0 (type: 5)
datistemplate = f (type: 1)
datallowconn = t (type: 1)
datconnlimit = -1 (type: 5)
datlastsysoid = 11510 (type: 1)
datfrozenxid = 379 (type: 1)
dattablespace = 1663 (type: 1)
datconfig = (type: 1)
datacl = (type: 1)
datid = 11511 (type: 1)
datname = postgres (type: 1)
numbackends = 0 (type: 5)
xact_commit = 221069 (type: 9)
xact_rollback = 18 (type: 9)
blks_read = 1176 (type: 9)
blks_hit = 13943750 (type: 9)
tup_returned = 77410091 (type: 9)
tup_fetched = 3253694 (type: 9)
tup_inserted = 0 (type: 9)
tup_updated = 0 (type: 9)
tup_deleted = 0 (type: 9)

36.8. Error Handling

This section describes how you can handle exceptional conditions and warnings in an embedded SQL program. There are two nonexclusive facilities for this.

- Callbacks can be configured to handle warning and error conditions using the WHENEVER command.
- Detailed information about the error or warning can be obtained from the sqlca variable.

36.8.1. Setting Callbacks

One simple method to catch errors and warnings is to set a specific action to be executed whenever a particular condition occurs. In general:

EXEC SQL WHENEVER condition action;

condition can be one of the following:

SQLERROR

The specified action is called whenever an error occurs during the execution of an SQL statement.

SQLWARNING

The specified action is called whenever a warning occurs during the execution of an SQL statement.

NOT FOUND

The specified action is called whenever an SQL statement retrieves or affects zero rows. (This condition is not an error, but you might be interested in handling it specially.)

action can be one of the following:

CONTINUE

This effectively means that the condition is ignored. This is the default.

GOTO label
GO TO label

Jump to the specified label (using a C goto statement).

SQLPRINT

Print a message to standard error. This is useful for simple programs or during prototyping. The details of the message cannot be configured.
STOP

Call exit(1), which will terminate the program.

DO BREAK

Execute the C statement break. This should only be used in loops or switch statements.

DO CONTINUE

Execute the C statement continue. This should only be used in loops statements. If executed, will cause the flow of control to return to the top of the loop.

CALL name (args)

DO name (args)

Call the specified C functions with the specified arguments. (This use is different from the meaning of CALL and DO in the normal PostgreSQL grammar.)

The SQL standard only provides for the actions CONTINUE and GOTO (and GO TO).

Here is an example that you might want to use in a simple program. It prints a simple message when a warning occurs and aborts the program when an error happens:

EXEC SQL WHENEVER SQLWARNING SQLPRINT;
EXEC SQL WHENEVER SQLERROR STOP;

The statement EXEC SQL WHENEVER is a directive of the SQL preprocessor, not a C statement. The error or warning actions that it sets apply to all embedded SQL statements that appear below the point where the handler is set, unless a different action was set for the same condition between the first EXEC SQL WHENEVER and the SQL statement causing the condition, regardless of the flow of control in the C program. So neither of the two following C program excerpts will have the desired effect:

```c
/*
 * WRONG
 */
int main(int argc, char *argv[])
{
    ...
    if (verbose) {
        EXEC SQL WHENEVER SQLWARNING SQLPRINT;
    }
    ...
    EXEC SQL SELECT ...;
    ...
}

/*
 * WRONG
 */
int main(int argc, char *argv[])
{
    ...
    set_error_handler();
    ...
    EXEC SQL SELECT ...;
```
...}

static void set_error_handler(void)
{
    EXEC SQL WHENEVER SQLERROR STOP;
}

36.8.2. sqlca

For more powerful error handling, the embedded SQL interface provides a global variable with the
name sqlca (SQL communication area) that has the following structure:

struct
{
    char sqlcaid[8];
    long sqlabc;
    long sqlcode;
    struct
    {
        int sqlerrml;
        char sqlerrmc[SQLERRMC_LEN];
    } sqlerrm;
    char sqlerrp[8];
    long sqlerrd[6];
    char sqlwarn[8];
    char sqlstate[5];
} sqlca;

(In a multithreaded program, every thread automatically gets its own copy of sqlca. This works
similarly to the handling of the standard C global variable errno.)

sqlca covers both warnings and errors. If multiple warnings or errors occur during the execution of
a statement, then sqlca will only contain information about the last one.

If no error occurred in the last SQL statement, sqlca.sqlcode will be 0 and sqlca.sqlstate
will be "00000". If a warning or error occurred, then sqlca.sqlcode will be negative and sql-
ca.sqlstate will be different from "00000". A positive sqlca.sqlcode indicates a harmless
condition, such as that the last query returned zero rows. sqlcode and sqlstate are two different
error code schemes; details appear below.

If the last SQL statement was successful, then sqlca.sqlerrd[1] contains the OID of the
processed row, if applicable, and sqlca.sqlerrd[2] contains the number of processed or returned
rows, if applicable to the command.

In case of an error or warning, sqlca.sqlerrm.sqlerrmc will contain a string that describes
the error. The field sqlca.sqlerm.sqlerrml contains the length of the error message that is
stored in sqlca.sqlerrm.sqlerrmc (the result of strlen(), not really interesting for a C
programmer). Note that some messages are too long to fit in the fixed-size sqlerrmc array; they
will be truncated.

In case of a warning, sqlca.sqlwarn[2] is set to W. (In all other cases, it is set to something
different from W.) If sqlca.sqlwarn[1] is set to W, then a value was truncated when it was stored
in a host variable. sqlca.sqlwarn[0] is set to W if any of the other elements are set to indicate
a warning.

The fields sqlcaid, sqlcabc, sqlerrp, and the remaining elements of sqlerrd and sqlwarn
currently contain no useful information.
The structure sqlca is not defined in the SQL standard, but is implemented in several other SQL database systems. The definitions are similar at the core, but if you want to write portable applications, then you should investigate the different implementations carefully.

Here is one example that combines the use of WHENEVER and sqlca, printing out the contents of sqlca when an error occurs. This is perhaps useful for debugging or prototyping applications, before installing a more “user-friendly” error handler.

EXEC SQL WHENEVER SQLERROR CALL print_sqlca();

void print_sqlca()
{
    fprintf(stderr, "==== sqlca ====
");
    fprintf(stderr, "sqlcode: %ld\n", sqlca.sqlcode);
    fprintf(stderr, "sqlerrm.sqlerrml: %d\n",
    sqlca.sqlerrm.sqlerrml);
    fprintf(stderr, "sqlerrm.sqlerrmc: %s\n",
    sqlca.sqlerrm.sqlerrmc);
    fprintf(stderr, "sqlerrd: %ld %ld %ld %ld %ld %ld\n",
    sqlca.sqlerrd[0],sqlca.sqlerrd[1],sqlca.sqlerrd[2],
    sqlca.sqlerrd[3],sqlca.sqlerrd[4],sqlca.sqlerrd[5]);
    fprintf(stderr, "sqlwarn: %d %d %d %d %d %d %d %d\n",
    sqlca.sqlwarn[0], sqlca.sqlwarn[1], sqlca.sqlwarn[2],
    sqlca.sqlwarn[3], sqlca.sqlwarn[4], sqlca.sqlwarn[5],
    sqlca.sqlwarn[6], sqlca.sqlwarn[7]);
    fprintf(stderr, "sqlstate: %5s\n", sqlca.sqlstate);
    fprintf(stderr, "===============\n");
}

The result could look as follows (here an error due to a misspelled table name):

==== sqlca ====
sqrcode: -400
sqlerrm.sqlerrml: 49
sqlerrm.sqlerrmc: relation "pg_databasep" does not exist on line 38
sqlerrd: 0 0 0 0 0 0
sqlwarn: 0 0 0 0 0 0 0 0
sqlstate: 42P01
===============

36.8.3. SQLSTATE VS. SQLCODE

The fields sqlca.sqlstate and sqlca.sqlcode are two different schemes that provide error codes. Both are derived from the SQL standard, but SQLCODE has been marked deprecated in the SQL-92 edition of the standard and has been dropped in later editions. Therefore, new applications are strongly encouraged to use SQLSTATE.

SQLSTATE is a five-character array. The five characters contain digits or upper-case letters that represent codes of various error and warning conditions. SQLSTATE has a hierarchical scheme: the first two characters indicate the general class of the condition, the last three characters indicate a subclass of the general condition. A successful state is indicated by the code 00000. The SQLSTATE codes are for the most part defined in the SQL standard. The PostgreSQL server natively supports SQLSTATE.
error codes; therefore a high degree of consistency can be achieved by using this error code scheme throughout all applications. For further information see Appendix A.

SQLCODE, the deprecated error code scheme, is a simple integer. A value of 0 indicates success, a positive value indicates success with additional information, a negative value indicates an error. The SQL standard only defines the positive value +100, which indicates that the last command returned or affected zero rows, and no specific negative values. Therefore, this scheme can only achieve poor portability and does not have a hierarchical code assignment. Historically, the embedded SQL processor for PostgreSQL has assigned some specific SQLCODE values for its use, which are listed below with their numeric value and their symbolic name. Remember that these are not portable to other SQL implementations. To simplify the porting of applications to the SQLSTATE scheme, the corresponding SQLSTATE is also listed. There is, however, no one-to-one or one-to-many mapping between the two schemes (indeed it is many-to-many), so you should consult the global SQLSTATE listing in Appendix A in each case.

These are the assigned SQLCODE values:

0 (ECPG_NO_ERROR)

Indicates no error. (SQLSTATE 00000)

100 (ECPG_NOT_FOUND)

This is a harmless condition indicating that the last command retrieved or processed zero rows, or that you are at the end of the cursor. (SQLSTATE 02000)

When processing a cursor in a loop, you could use this code as a way to detect when to abort the loop, like this:

```c
while (1)
{
    EXEC SQL FETCH ... ;
    if (sqlca.sqlcode == ECPG_NOT_FOUND)
        break;
}
```

But WHENEVER NOT FOUND DO BREAK effectively does this internally, so there is usually no advantage in writing this out explicitly.

-12 (ECPG_OUT_OF_MEMORY)

Indicates that your virtual memory is exhausted. The numeric value is defined as −ENOMEM. (SQLSTATE YE001)

-200 (ECPG_UNSUPPORTED)

Indicates the preprocessor has generated something that the library does not know about. Perhaps you are running incompatible versions of the preprocessor and the library. (SQLSTATE YE002)

-201 (ECPG_TOO_MANY_ARGUMENTS)

This means that the command specified more host variables than the command expected. (SQLSTATE 07001 or 07002)

-202 (ECPG_TOO_FEW_ARGUMENTS)

This means that the command specified fewer host variables than the command expected. (SQLSTATE 07001 or 07002)
This means a query has returned multiple rows but the statement was only prepared to store one result row (for example, because the specified variables are not arrays). (SQLSTATE 21000)

The host variable is of type int and the datum in the database is of a different type and contains a value that cannot be interpreted as an int. The library uses strtol() for this conversion. (SQLSTATE 42804)

The host variable is of type unsigned int and the datum in the database is of a different type and contains a value that cannot be interpreted as an unsigned int. The library uses strtoul() for this conversion. (SQLSTATE 42804)

The host variable is of type float and the datum in the database is of another type and contains a value that cannot be interpreted as a float. The library uses strtod() for this conversion. (SQLSTATE 42804)

The host variable is of type numeric and the datum in the database is of another type and contains a value that cannot be interpreted as a numeric value. (SQLSTATE 42804)

The host variable is of type interval and the datum in the database is of another type and contains a value that cannot be interpreted as an interval value. (SQLSTATE 42804)

The host variable is of type date and the datum in the database is of another type and contains a value that cannot be interpreted as a date value. (SQLSTATE 42804)

The host variable is of type timestamp and the datum in the database is of another type and contains a value that cannot be interpreted as a timestamp value. (SQLSTATE 42804)

This means the host variable is of type bool and the datum in the database is neither 't' nor 'f'. (SQLSTATE 42804)

The statement sent to the PostgreSQL server was empty. (This cannot normally happen in an embedded SQL program, so it might point to an internal error.) (SQLSTATE YE002)

A null value was returned and no null indicator variable was supplied. (SQLSTATE 22002)

An ordinary variable was used in a place that requires an array. (SQLSTATE 42804)
-215 (ECPG_DATA_NOT_ARRAY)

The database returned an ordinary variable in a place that requires array value. (SQLSTATE 42804)

-216 (ECPG_ARRAY_INSERT)

The value could not be inserted into the array. (SQLSTATE 42804)

-220 (ECPG_NO_CONN)

The program tried to access a connection that does not exist. (SQLSTATE 08003)

-221 (ECPG_NOT_CONN)

The program tried to access a connection that does exist but is not open. (This is an internal error.) (SQLSTATE YE002)

-230 (ECPG_INVALID_STMT)

The statement you are trying to use has not been prepared. (SQLSTATE 26000)

-239 (ECPG_INFORMIX_DUPLICATE_KEY)

Duplicate key error, violation of unique constraint (Informix compatibility mode). (SQLSTATE 23505)

-240 (ECPG_UNKNOWN_DESCRIPTOR)

The descriptor specified was not found. The statement you are trying to use has not been prepared. (SQLSTATE 33000)

-241 (ECPG_INVALID_DESCRIPTOR_INDEX)

The descriptor index specified was out of range. (SQLSTATE 07009)

-242 (ECPG_UNKNOWN_DESCRIPTOR_ITEM)

An invalid descriptor item was requested. (This is an internal error.) (SQLSTATE YE002)

-243 (ECPG_VAR_NOT_NUMERIC)

During the execution of a dynamic statement, the database returned a numeric value and the host variable was not numeric. (SQLSTATE 07006)

-244 (ECPG_VAR_NOT_CHAR)

During the execution of a dynamic statement, the database returned a non-numeric value and the host variable was numeric. (SQLSTATE 07006)

-284 (ECPG_INFORMIX_SUBSELECT_NOT_ONE)

A result of the subquery is not single row (Informix compatibility mode). (SQLSTATE 21000)

-400 (ECPG_PGSQ)

Some error caused by the PostgreSQL server. The message contains the error message from the PostgreSQL server.

-401 (ECPG_TRANS)

The PostgreSQL server signaled that we cannot start, commit, or rollback the transaction. (SQLSTATE 08007)
-402 (ECPG_CONNECT)

The connection attempt to the database did not succeed. (SQLSTATE 08001)

-403 (ECPG_DUPLICATE_KEY)

Duplicate key error, violation of unique constraint. (SQLSTATE 23505)

-404 (ECPG_SUBSELECT_NOT_ONE)

A result for the subquery is not single row. (SQLSTATE 21000)

-602 (ECPG_WARNING_UNSUPPORTED_PORTAL)

An invalid cursor name was specified. (SQLSTATE 34000)

-603 (ECPG_WARNING_IN_TRANSACTION)

Transaction is in progress. (SQLSTATE 25001)

-604 (ECPG_WARNING_NO_TRANSACTION)

There is no active (in-progress) transaction. (SQLSTATE 25P01)

-605 (ECPG_WARNING_PORTAL_EXISTS)

An existing cursor name was specified. (SQLSTATE 42P03)

### 36.9. Preprocessor Directives

Several preprocessor directives are available that modify how the _ecpg_ preprocessor parses and processes a file.

#### 36.9.1. Including Files

To include an external file into your embedded SQL program, use:

```sql
EXEC SQL INCLUDE filename;
EXEC SQL INCLUDE <filename>;
EXEC SQL INCLUDE "filename";
```

The embedded SQL preprocessor will look for a file named `filename`.h, preprocess it, and include it in the resulting C output. Thus, embedded SQL statements in the included file are handled correctly.

The _ecpg_ preprocessor will search a file at several directories in following order:

- current directory
- /usr/local/include
- PostgreSQL include directory, defined at build time (e.g., /usr/local/pgsql/include)
- /usr/include

But when EXEC SQL INCLUDE "filename" is used, only the current directory is searched.

In each directory, the preprocessor will first look for the file name as given, and if not found will append .h to the file name and try again (unless the specified file name already has that suffix).
Note that EXEC SQL INCLUDE is not the same as:

\#include <filename.h>

because this file would not be subject to SQL command preprocessing. Naturally, you can continue to use the C \#include directive to include other header files.

Note

The include file name is case-sensitive, even though the rest of the EXEC SQL INCLUDE command follows the normal SQL case-sensitivity rules.

36.9.2. The define and undef Directives

Similar to the directive \#define that is known from C, embedded SQL has a similar concept:

EXEC SQL DEFINE name;
EXEC SQL DEFINE name value;

So you can define a name:

EXEC SQL DEFINE HAVEFEATURE;

And you can also define constants:

EXEC SQL DEFINE MYNUMBER 12;
EXEC SQL DEFINE MYSTRING 'abc';

Use undef to remove a previous definition:

EXEC SQL UNDEF MYNUMBER;

Of course you can continue to use the C versions \#define and \#undef in your embedded SQL program. The difference is where your defined values get evaluated. If you use EXEC SQL DEFINE then the ecpg preprocessor evaluates the defines and substitutes the values. For example if you write:

EXEC SQL DEFINE MYNUMBER 12;
...
EXEC SQL UPDATE Tbl SET col = MYNUMBER;

then ecpg will already do the substitution and your C compiler will never see any name or identifier MYNUMBER. Note that you cannot use \#define for a constant that you are going to use in an embedded SQL query because in this case the embedded SQL precompiler is not able to see this declaration.

36.9.3. ifdef, ifndef, else, elif, and endif Directives

You can use the following directives to compile code sections conditionally:

EXEC SQL ifdef name;

Checks a name and processes subsequent lines if name has been created with EXEC SQL define name.
EXEC SQL ifndef name;

Checks a name and processes subsequent lines if name has not been created with EXEC SQL define name.

EXEC SQL else;

Starts processing an alternative section to a section introduced by either EXEC SQL ifdef name or EXEC SQL ifndef name.

EXEC SQL elif name;

Checks name and starts an alternative section if name has been created with EXEC SQL define name.

EXEC SQL endif;

Ends an alternative section.

Example:

EXEC SQL ifndef TZVAR;
EXEC SQL SET TIMEZONE TO 'GMT';
EXEC SQL elif TZNAME;
EXEC SQL SET TIMEZONE TO TZNAME;
EXEC SQL else;
EXEC SQL SET TIMEZONE TO TZVAR;
EXEC SQL endif;

36.10. Processing Embedded SQL Programs

Now that you have an idea how to form embedded SQL C programs, you probably want to know how to compile them. Before compiling you run the file through the embedded SQL C preprocessor, which converts the SQL statements you used to special function calls. After compiling, you must link with a special library that contains the needed functions. These functions fetch information from the arguments, perform the SQL command using the libpq interface, and put the result in the arguments specified for output.

The preprocessor program is called ecpg and is included in a normal PostgreSQL installation. Embedded SQL programs are typically named with an extension .pgc. If you have a program file called prog1.pgc, you can preprocess it by simply calling:

ecpg prog1.pgc

This will create a file called prog1.c. If your input files do not follow the suggested naming pattern, you can specify the output file explicitly using the -o option.

The preprocessed file can be compiled normally, for example:

c -c prog1.c

The generated C source files include header files from the PostgreSQL installation, so if you installed PostgreSQL in a location that is not searched by default, you have to add an option such as -I/usr/local/pgsql/include to the compilation command line.

To link an embedded SQL program, you need to include the libecpg library, like so:

c -o myprog prog1.o prog2.o ... -lecpg
Again, you might have to add an option like `-L/usr/local/pgsql/lib` to that command line.

You can use `pg_config` or `pkg-config` with package name `libecpg` to get the paths for your installation.

If you manage the build process of a larger project using `make`, it might be convenient to include the following implicit rule to your `makefiles`:

```makefile
ECPG = ecpg
%.c: %.pgc
   $(ECPG) $<
```

The complete syntax of the `ecpg` command is detailed in `ecpg`.

The `ecpg` library is thread-safe by default. However, you might need to use some threading command-line options to compile your client code.

### 36.11. Library Functions

The `libecpg` library primarily contains “hidden” functions that are used to implement the functionality expressed by the embedded SQL commands. But there are some functions that can usefully be called directly. Note that this makes your code unportable.

- `ECPGdebug(int on, FILE *stream)` turns on debug logging if called with the first argument non-zero. Debug logging is done on `stream`. The log contains all SQL statements with all the input variables inserted, and the results from the PostgreSQL server. This can be very useful when searching for errors in your SQL statements.

  **Note**

  On Windows, if the `ecpg` libraries and an application are compiled with different flags, this function call will crash the application because the internal representation of the `FILE` pointers differ. Specifically, multithreaded/single-threaded, release/debug, and static/dynamic flags should be the same for the library and all applications using that library.

- `ECPGget_PGconn(const char *connection_name)` returns the library database connection handle identified by the given name. If `connection_name` is set to `NULL`, the current connection handle is returned. If no connection handle can be identified, the function returns `NULL`. The returned connection handle can be used to call any other functions from `libpq`, if necessary.

  **Note**

  It is a bad idea to manipulate database connection handles made from `ecpg` directly with `libpq` routines.

- `ECPGtransactionStatus(const char *connection_name)` returns the current transaction status of the given connection identified by `connection_name`. See Section 34.2 and `libpq`'s `PQtransactionStatus()` for details about the returned status codes.

- `ECPGstatus(int lineno, const char* connection_name)` returns true if you are connected to a database and false if not. `connection_name` can be `NULL` if a single connection is being used.
**36.12. Large Objects**

Large objects are not directly supported by ECPG, but ECPG application can manipulate large objects through the libpq large object functions, obtaining the necessary PGconn object by calling the `ECPGget_PGconn()` function. (However, use of the `ECPGget_PGconn()` function and touching PGconn objects directly should be done very carefully and ideally not mixed with other ECPG database access calls.)

For more details about the `ECPGget_PGconn()`, see Section 36.11. For information about the large object function interface, see Chapter 35.

Large object functions have to be called in a transaction block, so when autocommit is off, `BEGIN` commands have to be issued explicitly.

Example 36.2 shows an example program that illustrates how to create, write, and read a large object in an ECPG application.

**Example 36.2. ECPG Program Accessing Large Objects**

```c
#include <stdio.h>
#include <stdlib.h>
#include <libpq-fe.h>
#include <libpq/libpq-fs.h>

EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
    PGconn *conn;
    Oid loid;
    int fd;
    char buf[256];
    int buflen = 256;
    char buf2[256];
    int rc;

    memset(buf, 1, buflen);

    EXEC SQL CONNECT TO testdb AS con1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

    conn = ECPGget_PGconn("con1");
    printf("conn = %p\n", conn);

    /* create */
    loid = lo_create(conn, 0);
    if (loid &lt; 0)
        printf("lo_create() failed: %s", PQerrorMessage(conn));

    printf("loid = %d\n", loid);

    /* write test */
    fd = lo_open(conn, loid, INV_READ|INV_WRITE);
    if (fd &lt; 0)
        printf("lo_open() failed: %s", PQerrorMessage(conn));
```

922
printf("fd = %d\n", fd);

rc = lo_write(conn, fd, buf, buflen);
if (rc &lt; 0)
    printf("lo_write() failed\n");

rc = lo_close(conn, fd);
if (rc &lt; 0)
    printf("lo_close() failed: %s", PQerrorMessage(conn));

/* read test */
fd = lo_open(conn, loid, INV_READ);
if (fd &lt; 0)
    printf("lo_open() failed: %s", PQerrorMessage(conn));

printf("fd = %d\n", fd);

rc = lo_read(conn, fd, buf2, buflen);
if (rc &lt; 0)
    printf("lo_read() failed\n");

rc = lo_close(conn, fd);
if (rc &lt; 0)
    printf("lo_close() failed: %s", PQerrorMessage(conn));

/* check */
rc = memcmp(buf, buf2, buflen);
printf("memcmp() = %d\n", rc);

/* cleanup */
rc = lo_unlink(conn, loid);
if (rc &lt; 0)
    printf("lo_unlink() failed: %s", PQerrorMessage(conn));

EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
return 0;
}

36.13. C++ Applications

ECPG has some limited support for C++ applications. This section describes some caveats.

The ecpg preprocessor takes an input file written in C (or something like C) and embedded SQL commands, converts the embedded SQL commands into C language chunks, and finally generates a .c file. The header file declarations of the library functions used by the C language chunks that ecpg generates are wrapped in extern "C" { ... } blocks when used under C++, so they should work seamlessly in C++.

In general, however, the ecpg preprocessor only understands C; it does not handle the special syntax and reserved words of the C++ language. So, some embedded SQL code written in C++ application code that uses complicated features specific to C++ might fail to be preprocessed correctly or might not work as expected.

A safe way to use the embedded SQL code in a C++ application is hiding the ECPG calls in a C module, which the C++ application code calls into to access the database, and linking that together with the rest of the C++ code. See Section 36.13.2 about that.
36.13.1. Scope for Host Variables

The ecpg preprocessor understands the scope of variables in C. In the C language, this is rather simple because the scopes of variables is based on their code blocks. In C++, however, the class member variables are referenced in a different code block from the declared position, so the ecpg preprocessor will not understand the scope of the class member variables.

For example, in the following case, the ecpg preprocessor cannot find any declaration for the variable dbname in the test method, so an error will occur.

```c++
class TestCpp
{
    EXEC SQL BEGIN DECLARE SECTION;
    char dbname[1024];
    EXEC SQL END DECLARE SECTION;

    public:
    TestCpp();
    void test();
    ~TestCpp();
};
TestCpp::TestCpp()
{
    EXEC SQL CONNECT TO testdb1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
}
void Test::test()
{
    EXEC SQL SELECT current_database() INTO :dbname;
    printf("current_database = %s\n", dbname);
}
TestCpp::~TestCpp()
{
    EXEC SQL DISCONNECT ALL;
}
```

This code will result in an error like this:

```bash
ecpg test_cpp.pgc
```

`test_cpp.pgc:28: ERROR: variable "dbname" is not declared`

To avoid this scope issue, the test method could be modified to use a local variable as intermediate storage. But this approach is only a poor workaround, because it uglifies the code and reduces performance.

```c++
void TestCpp::test()
{
    EXEC SQL BEGIN DECLARE SECTION;
    char tmp[1024];
    EXEC SQL END DECLARE SECTION;

    EXEC SQL SELECT current_database() INTO :tmp;
```
### 36.13.2. C++ Application Development with External C Module

If you understand these technical limitations of the `ecpg` preprocessor in C++, you might come to the conclusion that linking C objects and C++ objects at the link stage to enable C++ applications to use ECPG features could be better than writing some embedded SQL commands in C++ code directly. This section describes a way to separate some embedded SQL commands from C++ application code with a simple example. In this example, the application is implemented in C++, while C and ECPG is used to connect to the PostgreSQL server.

Three kinds of files have to be created: a C file (*.pgc), a header file, and a C++ file:

**test_mod.pgc**

A sub-routine module to execute SQL commands embedded in C. It is going to be converted into `test_mod.c` by the preprocessor.

```c
#include "test_mod.h"
#include <stdio.h>

void
db_connect()
{
    EXEC SQL CONNECT TO testdb1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
}

void
db_test()
{
    EXEC SQL BEGIN DECLARE SECTION;
    char dbname[1024];
    EXEC SQL END DECLARE SECTION;

    EXEC SQL SELECT current_database() INTO :dbname;
    printf("current_database = %s\n", dbname);
}

void
db_disconnect()
{
    EXEC SQL DISCONNECT ALL;
}
```

**test_mod.h**

A header file with declarations of the functions in the C module (`test_mod.pgc`). It is included by `test_cpp.cpp`. This file has to have an `extern "C"` block around the declarations, because it will be linked from the C++ module.

```c
#ifdef __cplusplus
```
extern "C" {
    #endif

    void db_connect();
    void db_test();
    void db_disconnect();

    #ifdef __cplusplus
    }
    #endif

test_cpp.cpp

    The main code for the application, including the main routine, and in this example a C++ class.

    #include "test_mod.h"

class TestCpp
{
    public:
        TestCpp();
        void test();
        ~TestCpp();
    }

    TestCpp::TestCpp()
    {
        db_connect();
    }

    void
    TestCpp::test()
    {
        db_test();
    }

    TestCpp::~TestCpp()
    {
        db_disconnect();
    }

    int
    main(void)
    {
        TestCpp *t = new TestCpp();

        t->test();
        return 0;
    }

    To build the application, proceed as follows. Convert test_mod.pgc into test_mod.c by running ecpg, and generate test_mod.o by compiling test_mod.c with the C compiler:

    ecpg -o test_mod.c test_mod.pgc
    cc -c test_mod.c -o test_mod.o

    Next, generate test_cpp.o by compiling test_cpp.cpp with the C++ compiler:
c++ -c test_cpp.cpp -o test_cpp.o

Finally, link these object files, test_cpp.o and test_mod.o, into one executable, using the C++ compiler driver:

c++ test_cpp.o test_mod.o -lecpg -o test_cpp

36.14. Embedded SQL Commands

This section describes all SQL commands that are specific to embedded SQL. Also refer to the SQL commands listed in SQL Commands, which can also be used in embedded SQL, unless stated otherwise.
ALLOCATE DESCRIPTOR

ALLOCATE DESCRIPTOR — allocate an SQL descriptor area

Synopsis

ALLOCATE DESCRIPTOR name

Description

ALLOCATE DESCRIPTOR allocates a new named SQL descriptor area, which can be used to exchange data between the PostgreSQL server and the host program.

Descriptor areas should be freed after use using the DEALLOCATE DESCRIPTOR command.

Parameters

name

A name of SQL descriptor, case sensitive. This can be an SQL identifier or a host variable.

Examples

EXEC SQL ALLOCATE DESCRIPTOR mydesc;

Compatibility

ALLOCATE DESCRIPTOR is specified in the SQL standard.

See Also

DEALLOCATE DESCRIPTOR, GET DESCRIPTOR, SET DESCRIPTOR
CONNECT

CONNECT — establish a database connection

Synopsis

CONNECT TO connection_target [ AS connection_name ]
[ USER connection_user ]
CONNECT TO DEFAULT
CONNECT connection_user
DATABASE connection_target

Description

The CONNECT command establishes a connection between the client and the PostgreSQL server.

Parameters

connection_target

connection_target specifies the target server of the connection on one of several forms.

[ database_name ][@host ][:port ]

Connect over TCP/IP

unix:postgresql://host[:port]/[database_name][?connection_option]

Connect over Unix-domain sockets

tcp:postgresql://host[:port]/[database_name][?connection_option]

Connect over TCP/IP

SQL string constant

containing a value in one of the above forms

host variable

host variable of type char[] or VARCHAR[] containing a value in one of the above forms

connection_object

An optional identifier for the connection, so that it can be referred to in other commands. This can be an SQL identifier or a host variable.

connection_user

The user name for the database connection.

This parameter can also specify user name and password, using one the forms

user_name/password, user_name IDENTIFIED BY password, or user_name USING password.

User name and password can be SQL identifiers, string constants, or host variables.
DEFAULT

Use all default connection parameters, as defined by libpq.

Examples

Here are several variants for specifying connection parameters:

EXEC SQL CONNECT TO "connectdb" AS main;
EXEC SQL CONNECT TO "connectdb" AS second;
EXEC SQL CONNECT TO "unix:postgresql://200.46.204.71/connectdb" AS main
  USER connectuser;
EXEC SQL CONNECT TO "unix:postgresql://localhost/connectdb" AS main
  USER connectuser;
EXEC SQL CONNECT TO 'connectdb' AS main;
EXEC SQL CONNECT TO 'unix:postgresql://localhost/connectdb' AS main
  USER :user;
EXEC SQL CONNECT TO :db AS :id;
EXEC SQL CONNECT TO :db USER connectuser USING :pw;
EXEC SQL CONNECT TO @localhost AS main USER connectdb;
EXEC SQL CONNECT TO AS main USER connectdb;
EXEC SQL CONNECT TO connectdb AS :id;
EXEC SQL CONNECT TO connectdb AS main USER connectuser/connectdb;
EXEC SQL CONNECT TO connectdb AS main;
EXEC SQL CONNECT TO connectdb@localhost AS main;
EXEC SQL CONNECT TO tcp:postgresql://localhost/ USER connectdb;
EXEC SQL CONNECT TO tcp:postgresql://localhost/connectdb USER
  connectuser IDENTIFIED BY connectpw;
EXEC SQL CONNECT TO tcp:postgresql://localhost:20/connectdb USER
  connectuser IDENTIFIED BY connectpw;
EXEC SQL CONNECT TO unix:postgresql://localhost/ AS main USER
  connectdb;
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb AS main
  USER connectuser;
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb USER
  connectuser IDENTIFIED BY "connectpw";
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb USER
  connectuser USING "connectpw";
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb?
  connect_timeout=14 USER connectuser;

Here is an example program that illustrates the use of host variables to specify connection parameters:

```c
int
main(void)
{
  EXEC SQL BEGIN DECLARE SECTION;
    char *dbname     = "testdb";    /* database name */
    char *user       = "testuser";  /* connection user name */
    char *connection = "tcp:postgresql://localhost:5432/testdb";
                                 /* connection string */
    char ver[256];                  /* buffer to store the version
                                  * string */
  EXEC SQL END DECLARE SECTION;

  ECPGdebug(1, stderr);
```

930
EXEC SQL CONNECT TO :dbname USER :user;
EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
EXEC SQL SELECT version() INTO :ver;
EXEC SQL DISCONNECT;
printf("version: %s\n", ver);
EXEC SQL CONNECT TO :connection USER :user;
EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
EXEC SQL SELECT version() INTO :ver;
EXEC SQL DISCONNECT;
printf("version: %s\n", ver);
return 0;
}

**Compatibility**

CONNECT is specified in the SQL standard, but the format of the connection parameters is implementation-specific.

**See Also**

DISCONNECT, SET CONNECTION
DEALLOCATE DESCRIPTOR

DEALLOCATE DESCRIPTOR — deallocate an SQL descriptor area

Synopsis

DEALLOCATE DESCRIPTOR name

Description

DEALLOCATE DESCRIPTOR deallocates a named SQL descriptor area.

Parameters

name

The name of the descriptor which is going to be deallocated. It is case sensitive. This can be an SQL identifier or a host variable.

Examples

EXEC SQL DEALLOCATE DESCRIPTOR mydesc;

Compatibility

DEALLOCATE DESCRIPTOR is specified in the SQL standard.

See Also

ALLOCATE DESCRIPTOR, GET DESCRIPTOR, SET DESCRIPTOR
DECLARE — define a cursor

Synopsis

DECLARE cursor_name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ]
CURSOR [ { WITH | WITHOUT } HOLD ] FOR prepared_name
DECLARE cursor_name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ]
CURSOR [ { WITH | WITHOUT } HOLD ] FOR query

Description

DECLARE declares a cursor for iterating over the result set of a prepared statement. This command has slightly different semantics from the direct SQL command DECLARE: Whereas the latter executes a query and prepares the result set for retrieval, this embedded SQL command merely declares a name as a “loop variable” for iterating over the result set of a query; the actual execution happens when the cursor is opened with the OPEN command.

Parameters

cursor_name

A cursor name, case sensitive. This can be an SQL identifier or a host variable.

prepared_name

The name of a prepared query, either as an SQL identifier or a host variable.

query

A SELECT or VALUES command which will provide the rows to be returned by the cursor.

For the meaning of the cursor options, see DECLARE.

Examples

Examples declaring a cursor for a query:

EXEC SQL DECLARE C CURSOR FOR SELECT * FROM My_Table;
EXEC SQL DECLARE C CURSOR FOR SELECT Item1 FROM T;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT version();

An example declaring a cursor for a prepared statement:

EXEC SQL PREPARE stmt1 AS SELECT version();
EXEC SQL DECLARE cur1 CURSOR FOR stmt1;

Compatibility

DECLARE is specified in the SQL standard.

See Also

OPEN, CLOSE, DECLARE
DESCRIBE

DESCRIBE — obtain information about a prepared statement or result set

Synopsis

DESCRIBE [ OUTPUT ] prepared_name USING [ SQL ]
  DESCRIPTOR descriptor_name
DESCRIBE [ OUTPUT ] prepared_name INTO [ SQL ]
  DESCRIPTOR descriptor_name
DESCRIBE [ OUTPUT ] prepared_name INTO sqlda_name

Description

DESCRIBE retrieves metadata information about the result columns contained in a prepared statement, without actually fetching a row.

Parameters

prepared_name

  The name of a prepared statement. This can be an SQL identifier or a host variable.

descriptor_name

  A descriptor name. It is case sensitive. It can be an SQL identifier or a host variable.

sqlda_name

  The name of an SQLDA variable.

Examples

EXEC SQL ALLOCATE DESCRIPTOR mydesc;
EXEC SQL PREPARE stmt1 FROM :sql_stmt;
EXEC SQL DESCRIBE stmt1 INTO SQL DESCRIPTOR mydesc;
EXEC SQL GET DESCRIPTOR mydesc VALUE 1 :charvar = NAME;
EXEC SQL DEALLOCATE DESCRIPTOR mydesc;

Compatibility

DESCRIBE is specified in the SQL standard.

See Also

ALLOCATE DESCRIPTOR, GET DESCRIPTOR
**DISCONNECT**

DISCONNECT — terminate a database connection

**Synopsis**

```c
DISCONNECT connection_name
DISCONNECT [ CURRENT ]
DISCONNECT DEFAULT
DISCONNECT ALL
```

**Description**

DISCONNECT closes a connection (or all connections) to the database.

**Parameters**

`connection_name`

A database connection name established by the `CONNECT` command.

`CURRENT`

Close the “current” connection, which is either the most recently opened connection, or the connection set by the `SET CONNECTION` command. This is also the default if no argument is given to the `DISCONNECT` command.

`DEFAULT`

Close the default connection.

`ALL`

Close all open connections.

**Examples**

```c
int main(void)
{
    EXEC SQL CONNECT TO testdb AS DEFAULT USER testuser;
    EXEC SQL CONNECT TO testdb AS con1 USER testuser;
    EXEC SQL CONNECT TO testdb AS con2 USER testuser;
    EXEC SQL CONNECT TO testdb AS con3 USER testuser;

    EXEC SQL DISCONNECT CURRENT;  /* close con3          */
    EXEC SQL DISCONNECT DEFAULT;  /* close DEFAULT       */
    EXEC SQL DISCONNECT ALL;      /* close con2 and con1 */

    return 0;
}
```

**Compatibility**

DISCONNECT is specified in the SQL standard.
See Also

CONNECT, SET CONNECTION
EXECUTE IMMEDIATE

EXECUTE IMMEDIATE — dynamically prepare and execute a statement

Synopsis

EXECUTE IMMEDIATE string

Description

EXECUTE IMMEDIATE immediately prepares and executes a dynamically specified SQL statement, without retrieving result rows.

Parameters

string

A literal C string or a host variable containing the SQL statement to be executed.

Examples

Here is an example that executes an INSERT statement using EXECUTE IMMEDIATE and a host variable named command:

```c
sprintf(command, "INSERT INTO test (name, amount, letter) VALUES ('db: 'r1''', 1, 'f')");
EXEC SQL EXECUTE IMMEDIATE :command;
```

Compatibility

EXECUTE IMMEDIATE is specified in the SQL standard.
GET DESCRIPTOR

GET DESCRIPTOR — get information from an SQL descriptor area

Synopsis

GET DESCRIPTOR descriptor_name :cvariable = descriptor_header_item
[, ... ]
GET DESCRIPTOR descriptor_name VALUE column_number :cvariable
= descriptor_item [, ... ]

Description

GET DESCRIPTOR retrieves information about a query result set from an SQL descriptor area and stores it into host variables. A descriptor area is typically populated using FETCH or SELECT before using this command to transfer the information into host language variables.

This command has two forms: The first form retrieves descriptor “header” items, which apply to the result set in its entirety. One example is the row count. The second form, which requires the column number as additional parameter, retrieves information about a particular column. Examples are the column name and the actual column value.

Parameters

descriptor_name

A descriptor name.

descriptor_header_item

A token identifying which header information item to retrieve. Only COUNT, to get the number of columns in the result set, is currently supported.

column_number

The number of the column about which information is to be retrieved. The count starts at 1.

descriptor_item

A token identifying which item of information about a column to retrieve. See Section 36.7.1 for a list of supported items.

cvariable

A host variable that will receive the data retrieved from the descriptor area.

Examples

An example to retrieve the number of columns in a result set:

EXEC SQL GET DESCRIPTOR d :d_count = COUNT;

An example to retrieve a data length in the first column:

EXEC SQL GET DESCRIPTOR d VALUE 1 :d_returned_octet_length = RETURNED_OCTET_LENGTH;
An example to retrieve the data body of the second column as a string:

EXEC SQL GET DESCRIPTOR d VALUE 2 :d_data = DATA;

Here is an example for a whole procedure of executing SELECT current_database(); and showing the number of columns, the column data length, and the column data:

```c
int
main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    int  d_count;
    char d_data[1024];
    int  d_returned_octet_length;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO testdb AS con1 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', 'false); EXEC SQL COMMIT;
    EXEC SQL ALLOCATE DESCRIPTOR d;

    /* Declare, open a cursor, and assign a descriptor to the cursor */
    EXEC SQL DECLARE cur CURSOR FOR SELECT current_database();
    EXEC SQL OPEN cur;
    EXEC SQL FETCH NEXT FROM cur INTO SQL DESCRIPTOR d;

    /* Get a number of total columns */
    EXEC SQL GET DESCRIPTOR d :d_count = COUNT;
    printf("d_count = %d\n", d_count);

    /* Get length of a returned column */
    EXEC SQL GET DESCRIPTOR d VALUE 1 :d_returned_octet_length = RETURNED_OCTET_LENGTH;
    printf("d_returned_octet_length = %d\n", d_returned_octet_length);

    /* Fetch the returned column as a string */
    EXEC SQL GET DESCRIPTOR d VALUE 1 :d_data = DATA;
    printf("d_data = %s\n", d_data);

    /* Closing */
    EXEC SQL CLOSE cur;
    EXEC SQL COMMIT;

    EXEC SQL DEALLOCATE DESCRIPTOR d;
    EXEC SQL DISCONNECT ALL;

    return 0;
}
```

When the example is executed, the result will look like this:

```
d_count = 1
d_returned_octet_length = 6
d_data = testdb
```
Compatibility

GET DESCRIPTOR is specified in the SQL standard.

See Also

ALLOCATE DESCRIPTOR, SET DESCRIPTOR
OPEN

OPEN — open a dynamic cursor

Synopsis

OPEN cursor_name
OPEN cursor_name USING value [, ... ]
OPEN cursor_name USING SQL DESCRIPTOR descriptor_name

Description

OPEN opens a cursor and optionally binds actual values to the placeholders in the cursor's declaration. The cursor must previously have been declared with the DECLARE command. The execution of OPEN causes the query to start executing on the server.

Parameters

cursor_name

The name of the cursor to be opened. This can be an SQL identifier or a host variable.

value

A value to be bound to a placeholder in the cursor. This can be an SQL constant, a host variable, or a host variable with indicator.

descriptor_name

The name of a descriptor containing values to be bound to the placeholders in the cursor. This can be an SQL identifier or a host variable.

Examples

EXEC SQL OPEN a;
EXEC SQL OPEN d USING 1, 'test';
EXEC SQL OPEN c1 USING SQL DESCRIPTOR mydesc;
EXEC SQL OPEN :curname1;

Compatibility

OPEN is specified in the SQL standard.

See Also

DECLARE, CLOSE
PREPARE

PREPARE — prepare a statement for execution

Synopsis

PREPARE name FROM string

Description

PREPARE prepares a statement dynamically specified as a string for execution. This is different from the direct SQL statement PREPARE, which can also be used in embedded programs. The EXECUTE command is used to execute either kind of prepared statement.

Parameters

prepared_name

An identifier for the prepared query.

string

A literal C string or a host variable containing a preparable statement, one of the SELECT, INSERT, UPDATE, or DELETE.

Examples

char *stmt = "SELECT * FROM test1 WHERE a = ? AND b = ?";
EXEC SQL ALLOCATE DESCRIPTOR outdesc;
EXEC SQL PREPARE foo FROM :stmt;
EXEC SQL EXECUTE foo USING SQL DESCRIPTOR indesc INTO SQL DESCRIPTOR outdesc;

Compatibility

PREPARE is specified in the SQL standard.

See Also

EXECUTE
SET AUTOCOMMIT

SET AUTOCOMMIT — set the autocommit behavior of the current session

Synopsis

SET AUTOCOMMIT { = | TO } { ON | OFF }

Description

SET AUTOCOMMIT sets the autocommit behavior of the current database session. By default, embedded SQL programs are not in autocommit mode, so COMMIT needs to be issued explicitly when desired. This command can change the session to autocommit mode, where each individual statement is committed implicitly.

Compatibility

SET AUTOCOMMIT is an extension of PostgreSQL ECPG.
SET CONNECTION

SET CONNECTION — select a database connection

Synopsis

```
SET CONNECTION [ TO | = ] connection_name
```

Description

SET CONNECTION sets the “current” database connection, which is the one that all commands use unless overridden.

Parameters

`connection_name`

A database connection name established by the CONNECT command.

`DEFAULT`

Set the connection to the default connection.

Examples

```
EXEC SQL SET CONNECTION TO con2;
EXEC SQL SET CONNECTION = con1;
```

Compatibility

SET CONNECTION is specified in the SQL standard.

See Also

CONNECT, DISCONNECT
**SET DESCRIPTOR**

SET DESCRIPTOR — set information in an SQL descriptor area

**Synopsis**

```sql
SET DESCRIPTOR descriptor_name descriptor_header_item = value
[, ... ]
SET DESCRIPTOR descriptor_name VALUE number descriptor_item = value
[, ...
```

**Description**

SET DESCRIPTOR populates an SQL descriptor area with values. The descriptor area is then typically used to bind parameters in a prepared query execution.

This command has two forms: The first form applies to the descriptor “header”, which is independent of a particular datum. The second form assigns values to particular datums, identified by number.

**Parameters**

- `descriptor_name`  
  A descriptor name.

- `descriptor_header_item`  
  A token identifying which header information item to set. Only `COUNT`, to set the number of descriptor items, is currently supported.

- `number`  
  The number of the descriptor item to set. The count starts at 1.

- `descriptor_item`  
  A token identifying which item of information to set in the descriptor. See Section 36.7.1 for a list of supported items.

- `value`  
  A value to store into the descriptor item. This can be an SQL constant or a host variable.

**Examples**

```sql
EXEC SQL SET DESCRIPTOR indesc COUNT = 1;
EXEC SQL SET DESCRIPTOR indesc VALUE 1 DATA = 2;
EXEC SQL SET DESCRIPTOR indesc VALUE 1 DATA = :val1;
EXEC SQL SET DESCRIPTOR indesc VALUE 2 INDICATOR = :val1, DATA = 'some string';
EXEC SQL SET DESCRIPTOR indesc VALUE 2 INDICATOR = :val2null, DATA = :val2;
```

**Compatibility**

SET DESCRIPTOR is specified in the SQL standard.
See Also

ALLOCATE DESCRIPTOR, GET DESCRIPTOR
TYPE

TYPE — define a new data type

Synopsis

TYPE type_name IS ctype

Description

The TYPE command defines a new C type. It is equivalent to putting a typedef into a declare section.

This command is only recognized when ecpg is run with the -c option.

Parameters

  type_name

  The name for the new type. It must be a valid C type name.

  ctype

  A C type specification.

Examples

EXEC SQL TYPE customer IS
  struct
  {  
    varchar name[50];
    int     phone;
  };

EXEC SQL TYPE cust_ind IS
  struct ind
  {  
    short   name_ind;
    short   phone_ind;
  };

EXEC SQL TYPE c IS char reference;
EXEC SQL TYPE ind IS union { int integer; short smallint; };
EXEC SQL TYPE intarray IS int[AMOUNT];
EXEC SQL TYPE str IS varchar[BUFFERSZ];
EXEC SQL TYPE string IS char[11];

Here is an example program that uses EXEC SQL TYPE:

EXEC SQL WHENEVER SQLERROR SQLPRINT;

EXEC SQL TYPE tt IS
  struct
  {  
    varchar v[256];
    int     i;
EXEC SQL TYPE tt_ind IS
   struct ind {
       short v_ind;
       short i_ind;
   };

int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    tt t;
    tt_ind t_ind;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO testdb AS con1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
    EXEC SQL SELECT current_database(), 256 INTO :t:t_ind LIMIT 1;

    printf("t.v = %s\n", t.v.arr);
    printf("t.i = %d\n", t.i);

    printf("t_ind.v_ind = %d\n", t_ind.v_ind);
    printf("t_ind.i_ind = %d\n", t_ind.i_ind);

    EXEC SQL DISCONNECT con1;
    return 0;
}

The output from this program looks like this:

t.v = testdb
t.i = 256
t_ind.v_ind = 0
t_ind.i_ind = 0

Compatibility

The TYPE command is a PostgreSQL extension.
VAR

VAR — define a variable

Synopsis

VAR varname IS ctype

Description

The VAR command assigns a new C data type to a host variable. The host variable must be previously declared in a declare section.

Parameters

varname

A C variable name.

cype

A C type specification.

Examples

Exec sql begin declare section;
short a;
exec sql end declare section;
EXEC SQL VAR a IS int;

Compatibility

The VAR command is a PostgreSQL extension.
WHENEVER

WHENEVER — specify the action to be taken when an SQL statement causes a specific class condition to be raised

Synopsis

WHENEVER { NOT FOUND | SQLERROR | SQLWARNING } action

Description

Define a behavior which is called on the special cases (Rows not found, SQL warnings or errors) in the result of SQL execution.

Parameters

See Section 36.8.1 for a description of the parameters.

Examples

EXEC SQL WHENEVER NOT FOUND CONTINUE;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
EXEC SQL WHENEVER NOT FOUND DO CONTINUE;
EXEC SQL WHENEVER SQLWARNING SQLPRINT;
EXEC SQL WHENEVER SQLWARNING DO warn();
EXEC SQL WHENEVER SQLERROR sqlprint;
EXEC SQL WHENEVER SQLERROR CALL print2();
EXEC SQL WHENEVER SQLERROR DO handle_error("select");
EXEC SQL WHENEVER SQLERROR DO sqlnotice(NULL, NONO);
EXEC SQL WHENEVER SQLERROR DO sqlprint();
EXEC SQL WHENEVER SQLERROR GOTO error_label;
EXEC SQL WHENEVER SQLERROR STOP;

A typical application is the use of WHENEVER NOT FOUND BREAK to handle looping through result sets:

int main(void)
{
    EXEC SQL CONNECT TO testdb AS con1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', 'hoge', false); EXEC SQL COMMIT;
    EXEC SQL ALLOCATE DESCRIPTOR d;
    EXEC SQL DECLARE cur CURSOR FOR SELECT current_database(), 'hoge', 256;
    EXEC SQL OPEN cur;

    /* when end of result set reached, break out of while loop */
    EXEC SQL WHENEVER NOT FOUND DO BREAK;

    while (1)
    {
        EXEC SQL FETCH NEXT FROM cur INTO SQL DESCRIPTOR d;
        ...
    }
}
Compatibility

WHENEVER is specified in the SQL standard, but most of the actions are PostgreSQL extensions.

36.15. Informix Compatibility Mode

epg can be run in a so-called Informix compatibility mode. If this mode is active, it tries to behave as if it were the Informix precompiler for Informix E/SQL. Generally spoken this will allow you to use the dollar sign instead of the EXEC SQL primitive to introduce embedded SQL commands:

```c
int j = 3;
CONNECT TO :dbname;
CREATE TABLE test(i INT PRIMARY KEY, j INT);
INSERT INTO test(i, j) VALUES (7, :j);
COMMIT;
```

There must not be any white space between the $ and a following preprocessor directive, that is, include, define, ifdef, etc. Otherwise, the preprocessor will parse the token as a host variable.

There are two compatibility modes: INFORMIX, INFORMIX_SE

When linking programs that use this compatibility mode, remember to link against libcompat that is shipped with ECPG.

Besides the previously explained syntactic sugar, the Informix compatibility mode ports some functions for input, output and transformation of data as well as embedded SQL statements known from E/SQL to ECPG.

Informix compatibility mode is closely connected to the pgtypeslib library of ECPG. pgtypeslib maps SQL data types to data types within the C host program and most of the additional functions of the Informix compatibility mode allow you to operate on those C host program types. Note however that the extent of the compatibility is limited. It does not try to copy Informix behavior; it allows you to do more or less the same operations and gives you functions that have the same name and the same basic behavior but it is no drop-in replacement if you are using Informix at the moment. Moreover, some of the data types are different. For example, PostgreSQL's datetime and interval types do not know about ranges like for example YEAR TO MINUTE so you won't find support in ECPG for that either.

36.15.1. Additional Types

The Informix-special "string" pseudo-type for storing right-trimmed character string data is now supported in Informix-mode without using typedef. In fact, in Informix-mode, ECPG refuses to process source files that contain typedef sometype string;
EXEC SQL BEGIN DECLARE SECTION;
string userid; /* this variable will contain trimmed data */
EXEC SQL END DECLARE SECTION;

EXEC SQL FETCH MYCUR INTO :userid;

36.15.2. Additional/Missing Embedded SQL Statements

CLOSE DATABASE

This statement closes the current connection. In fact, this is a synonym for ECPG’s DISCONNECT
CURRENT:

$CLOSE DATABASE;                /* close the current connection
 * /
EXEC SQL CLOSE DATABASE;

FREE cursor_name

Due to the differences how ECPG works compared to Informix's ESQL/C (i.e. which steps are
purely grammar transformations and which steps rely on the underlying run-time library) there is
no FREE cursor_name statement in ECPG. This is because in ECPG, DECLARE CURSOR
doesn't translate to a function call into the run-time library that uses to the cursor name. This
means that there's no run-time bookkeeping of SQL cursors in the ECPG run-time library, only
in the PostgreSQL server.

FREE statement_name

FREE statement_name is a synonym for DEALLOCATE PREPARE statement_name.

36.15.3. Informix-compatible SQLDA Descriptor Areas

Informix-compatible mode supports a different structure than the one described in Section 36.7.2. See
below:

struct sqlvar_compat
{
    short sqltype;
    int sqllen;
    char *sqldata;
    short *sqlind;
    char *sqlname;
    char *sqlformat;
    short sqltype;
    short sqllilen;
    char *sqlidata;
    int sqlxid;
    char *sqltypename;
    short sqltypelen;
    short sqlownerlen;
    short sqlsourcetype;
    char *sqlownername;
    int sqlsourceid;
    char *sqlilongdata;
int sqlflags;
void *sqlreserved;
};

struct sqlda_compat
{
    short sqld;
    struct sqlvar_compat *sqlvar;
    char desc_name[19];
    short desc_occ;
    struct sqlda_compat *desc_next;
    void *reserved;
};

typedef struct sqlvar_compat sqlvar_t;
typedef struct sqlda_compat sqlda_t;

The global properties are:

sqlld

The number of fields in the SQLDA descriptor.

sqlvar

Pointer to the per-field properties.

desc_name

Unused, filled with zero-bytes.

desc_occ

Size of the allocated structure.

desc_next

Pointer to the next SQLDA structure if the result set contains more than one record.

reserved

Unused pointer, contains NULL. Kept for Informix-compatibility.

The per-field properties are below, they are stored in the sqlvar array:

sqltype

Type of the field. Constants are in sqltypes.h

sqlllen

Length of the field data.

sqlldata

Pointer to the field data. The pointer is of char * type, the data pointed by it is in a binary format. Example:

    int intval;
    switch (sqlldata->sqlvar[i].sqltype)
ECPG - Embedded SQL in C

{  
  case SQLINTEGER:
    intval = *(int *)sqldata->sqlvar[i].sqldata;
    break;
...  
}

sqlind

Pointer to the NULL indicator. If returned by DESCRIBE or FETCH then it's always a valid pointer. If used as input for EXECUTE ... USING sqlda; then NULL-pointer value means that the value for this field is non-NULL. Otherwise a valid pointer and sqlitype has to be properly set. Example:

if (*(int2 *)sqldata->sqlvar[i].sqlind != 0)
  printf("value is NULL\n");

sqlname

Name of the field. 0-terminated string.

sqlformat

Reserved in Informix, value of PQfformat() for the field.

sqlitype

Type of the NULL indicator data. It's always SQLSMINT when returning data from the server. When the SQLDA is used for a parameterized query, the data is treated according to the set type.

sqlilen

Length of the NULL indicator data.

sqlxid

Extended type of the field, result of PQftype().

sqltypename

sqltypelen

sqlownerlen

sqlsourcetype

sqlownernname

sqlsourceid

sqlflags

sqlreserved

Unused.

sqlilongdata

It equals to sqldata if sqllen is larger than 32kB.

Example:

EXEC SQL INCLUDE sqlda.h;

    sqlda_t      *sqlda; /* This doesn't need to be under embedded DECLARE SECTION */
EXEC SQL BEGIN DECLARE SECTION;
char *prep_stmt = "select * from table1";
int i;
EXEC SQL END DECLARE SECTION;
...
EXEC SQL PREPARE mystmt FROM :prep_stmt;
EXEC SQL DESCRIBE mystmt INTO sqlda;
printf("# of fields: %d\n", sqlda->sqld);
for (i = 0; i < sqlda->sqld; i++)
    printf("field %d: \"%s\"\n", sqlda->sqlvar[i]->sqlname);
EXEC SQL DECLARE mycursor CURSOR FOR mystmt;
EXEC SQL OPEN mycursor;
EXEC SQL WHENEVER NOT FOUND GOTO out;
while (1)
{
    EXEC SQL FETCH mycursor USING sqlda;
}
EXEC SQL CLOSE mycursor;
free(sqlda); /* The main structure is all to be free()
* sqlda and sqlda->sqlvar is in one allocated
area */

For more information, see the sqlda.h header and the src/interfaces/ecpg/test/com-
pat_informix/sqlda.pgc regression test.

### 36.15.4. Additional Functions

decadd

Add two decimal type values.

```c
int decadd(decimal *arg1, decimal *arg2, decimal *sum);
```

The function receives a pointer to the first operand of type decimal (arg1), a pointer to the second operand of type decimal (arg2) and a pointer to a value of type decimal that will contain the sum (sum). On success, the function returns 0. ECPG_INFORMIX_NUM_OVERFLOW is returned in case of overflow and ECPG_INFORMIX_NUM_UNDERFLOW in case of underflow. -1 is returned for other failures and errno is set to the respective errno number of the pgtypeslib.

deccmp

Compare two variables of type decimal.

```c
int deccmp(decimal *arg1, decimal *arg2);
```

The function receives a pointer to the first decimal value (arg1), a pointer to the second decimal value (arg2) and returns an integer value that indicates which is the bigger value.

- 1, if the value that arg1 points to is bigger than the value that var2 points to
• -1, if the value that \texttt{arg1} points to is smaller than the value that \texttt{arg2} points to
• 0, if the value that \texttt{arg1} points to and the value that \texttt{arg2} points to are equal

deccopy
Copy a decimal value.

\begin{verbatim}
void deccopy(decimal *src, decimal *target);
\end{verbatim}

The function receives a pointer to the decimal value that should be copied as the first argument \texttt{(src)} and a pointer to the target structure of type decimal \texttt{(target)} as the second argument.

deccvasc
Convert a value from its ASCII representation into a decimal type.

\begin{verbatim}
int deccvasc(char *cp, int len, decimal *np);
\end{verbatim}

The function receives a pointer to string that contains the string representation of the number to be converted \texttt{(cp)} as well as its length \texttt{len}. \texttt{np} is a pointer to the decimal value that saves the result of the operation.

Valid formats are for example: \texttt{-2.794}, \texttt{+3.44}, \texttt{592.49E07} or \texttt{-32.84e-4}.

The function returns 0 on success. If overflow or underflow occurred, \texttt{ECPG_INFORMIX_NUM_OVERFLOW} or \texttt{ECPG_INFORMIX_NUM_UNDERFLOW} is returned. If the ASCII representation could not be parsed, \texttt{ECPG_INFORMIX_BAD_NUMERIC} is returned or \texttt{ECPG_INFORMIX_BAD_EXPONENT} if this problem occurred while parsing the exponent.

deccvdbl
Convert a value of type double to a value of type decimal.

\begin{verbatim}
int deccvdbl(double dbl, decimal *np);
\end{verbatim}

The function receives the variable of type double that should be converted as its first argument \texttt{(dbl)}. As the second argument \texttt{(np)}, the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

deccvint
Convert a value of type int to a value of type decimal.

\begin{verbatim}
int deccvint(int in, decimal *np);
\end{verbatim}

The function receives the variable of type int that should be converted as its first argument \texttt{(in)}. As the second argument \texttt{(np)}, the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

deccvlong
Convert a value of type long to a value of type decimal.
int deccvlong(long lng, decimal *np);

The function receives the variable of type long that should be converted as its first argument (lng). As the second argument (np), the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

decdiv

Divide two variables of type decimal.

int decdiv(decimal *n1, decimal *n2, decimal *result);

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates n1/n2. result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the division fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively. If an attempt to divide by zero is observed, the function returns ECPG_INFORMIX_DIVIDE_ZERO.

decmul

Multiply two decimal values.

int decmul(decimal *n1, decimal *n2, decimal *result);

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates n1*n2. result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the multiplication fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively.

decsub

Subtract one decimal value from another.

int decsub(decimal *n1, decimal *n2, decimal *result);

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates n1-n2. result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the subtraction fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively.

dectoasc

Convert a variable of type decimal to its ASCII representation in a C char* string.

int dectoasc(decimal *np, char *cp, int len, int right)
The function receives a pointer to a variable of type decimal (\(np\)) that it converts to its textual representation. \(cp\) is the buffer that should hold the result of the operation. The parameter right specifies, how many digits right of the decimal point should be included in the output. The result will be rounded to this number of decimal digits. Setting right to -1 indicates that all available decimal digits should be included in the output. If the length of the output buffer, which is indicated by len is not sufficient to hold the textual representation including the trailing zero byte, only a single * character is stored in the result and -1 is returned.

The function returns either -1 if the buffer \(cp\) was too small or ECPG_INFORMIX_OUT_OF_MEMORY if memory was exhausted.

\[\text{dectodbl}\]

Convert a variable of type decimal to a double.

\[
\text{int dectodbl(decimal *np, double *dblp)};
\]

The function receives a pointer to the decimal value to convert (\(np\)) and a pointer to the double variable that should hold the result of the operation (\(dblp\)).

On success, 0 is returned and a negative value if the conversion failed.

\[\text{dectoint}\]

Convert a variable to type decimal to an integer.

\[
\text{int dectoint(decimal *np, int *ip)};
\]

The function receives a pointer to the decimal value to convert (\(np\)) and a pointer to the integer variable that should hold the result of the operation (\(ip\)).

On success, 0 is returned and a negative value if the conversion failed. If an overflow occurred, ECPG_INFORMIX_NUM_OVERFLOW is returned.

Note that the ECPG implementation differs from the Informix implementation. Informix limits an integer to the range from -32767 to 32767, while the limits in the ECPG implementation depend on the architecture (-INT_MAX .. INT_MAX).

\[\text{dectolong}\]

Convert a variable to type decimal to a long integer.

\[
\text{int dectolong(decimal *np, long *lngp)};
\]

The function receives a pointer to the decimal value to convert (\(np\)) and a pointer to the long variable that should hold the result of the operation (\(lngp\)).

On success, 0 is returned and a negative value if the conversion failed. If an overflow occurred, ECPG_INFORMIX_NUM_OVERFLOW is returned.

Note that the ECPG implementation differs from the Informix implementation. Informix limits a long integer to the range from -2,147,483,647 to 2,147,483,647, while the limits in the ECPG implementation depend on the architecture (-LONG_MAX .. LONG_MAX).

\[\text{rdatestr}\]

Converts a date to a C char* string.
int rdatestr(date d, char *str);

The function receives two arguments, the first one is the date to convert (d) and the second one is a pointer to the target string. The output format is always yyyy-mm-dd, so you need to allocate at least 11 bytes (including the zero-byte terminator) for the string.

The function returns 0 on success and a negative value in case of error.

Note that ECPG’s implementation differs from the Informix implementation. In Informix the format can be influenced by setting environment variables. In ECPG however, you cannot change the output format.

rstrdate

Parse the textual representation of a date.

int rstrdate(char *str, date *d);

The function receives the textual representation of the date to convert (str) and a pointer to a variable of type date (d). This function does not allow you to specify a format mask. It uses the default format mask of Informix which is mm/dd/yyyy. Internally, this function is implemented by means of rdefmtdate. Therefore, rstrdate is not faster and if you have the choice you should opt for rdefmtdate which allows you to specify the format mask explicitly.

The function returns the same values as rdefmtdate.

rtoday

Get the current date.

void rtoday(date *d);

The function receives a pointer to a date variable (d) that it sets to the current date.

Internally this function uses the PGTYPESdate_today function.

rjulmdy

Extract the values for the day, the month and the year from a variable of type date.

int rjulmdy(date d, short mdy[3]);

The function receives the date d and a pointer to an array of 3 short integer values mdy. The variable name indicates the sequential order: mdy[0] will be set to contain the number of the month, mdy[1] will be set to the value of the day and mdy[2] will contain the year.

The function always returns 0 at the moment.

Internally the function uses the PGTYPESdate_julmdy function.

rdefmtdate

Use a format mask to convert a character string to a value of type date.

int rdefmtdate(date *d, char *fmt, char *str);

The function receives a pointer to the date value that should hold the result of the operation (d), the format mask to use for parsing the date (fmt) and the C char* string containing the textual
representation of the date (str). The textual representation is expected to match the format mask. However you do not need to have a 1:1 mapping of the string to the format mask. The function only analyzes the sequential order and looks for the literals yy or yyyy that indicate the position of the year, mm to indicate the position of the month and dd to indicate the position of the day.

The function returns the following values:

- 0 - The function terminated successfully.
- ECPG_INFORMIX_ENOSHORTDATE - The date does not contain delimiters between day, month and year. In this case the input string must be exactly 6 or 8 bytes long but isn't.
- ECPG_INFORMIX_ENOTDMY - The format string did not correctly indicate the sequential order of year, month and day.
- ECPG_INFORMIX_BAD_DAY - The input string does not contain a valid day.
- ECPG_INFORMIX_BAD_MONTH - The input string does not contain a valid month.
- ECPG_INFORMIX_BAD_YEAR - The input string does not contain a valid year.

Internally this function is implemented to use the PGTYPESdate_defmt_asc function. See the reference there for a table of example input.

**rfmtdate**

Convert a variable of type date to its textual representation using a format mask.

```
int rfmtdate(date d, char *fmt, char *str);
```

The function receives the date to convert (d), the format mask (fmt) and the string that will hold the textual representation of the date (str).

On success, 0 is returned and a negative value if an error occurred.

Internally this function uses the PGTYPESdate_fmt_asc function, see the reference there for examples.

**rmdyjul**

Create a date value from an array of 3 short integers that specify the day, the month and the year of the date.

```
int rmdyjul(short mdy[3], date *d);
```

The function receives the array of the 3 short integers (mdy) and a pointer to a variable of type date that should hold the result of the operation.

Currently the function returns always 0.

Internally the function is implemented to use the function PGTYPESdate_mdyjul.

**rdayofweek**

Return a number representing the day of the week for a date value.

```
int rdayofweek(date d);
```

The function receives the date variable d as its only argument and returns an integer that indicates the day of the week for this date.
Internally the function is implemented to use the function PGTYPES\texttt{date\_dayofweek}.

\texttt{dtcurrent}

Retrieve the current timestamp.

\begin{verbatim}
void dtcurrent(timestamp *ts);
\end{verbatim}

The function retrieves the current timestamp and saves it into the timestamp variable that \texttt{ts} points to.

\texttt{dtcvasc}

Parses a timestamp from its textual representation into a timestamp variable.

\begin{verbatim}
int dtcvasc(char *str, timestamp *ts);
\end{verbatim}

The function receives the string to parse (\texttt{str}) and a pointer to the timestamp variable that should hold the result of the operation (\texttt{ts}).

The function returns 0 on success and a negative value in case of error.

Internally this function uses the PGTYPE\texttt{Stimestamp\_from\_asc} function. See the reference there for a table with example inputs.

\texttt{dtcvfmtasc}

Parses a timestamp from its textual representation using a format mask into a timestamp variable.

\begin{verbatim}
dtcvfmtasc(char *inbuf, char *fmtstr, timestamp *dtvalue)
\end{verbatim}

The function receives the string to parse (\texttt{inbuf}), the format mask to use (\texttt{fmtstr}) and a pointer to the timestamp variable that should hold the result of the operation (\texttt{dtvalue}).

This function is implemented by means of the PGTYPE\texttt{Stimestamp\_defmt\_asc} function. See the documentation there for a list of format specifiers that can be used.

The function returns 0 on success and a negative value in case of error.

\texttt{dtsub}

Subtract one timestamp from another and return a variable of type interval.

\begin{verbatim}
int dtsub(timestamp *ts1, timestamp *ts2, interval *iv);
\end{verbatim}
The function will subtract the timestamp variable that `ts2` points to from the timestamp variable that `ts1` points to and will store the result in the interval variable that `iv` points to.

Upon success, the function returns 0 and a negative value if an error occurred.

dttoasc

Convert a timestamp variable to a C char* string.

```c
int dttoasc(timestamp *ts, char *output);
```

The function receives a pointer to the timestamp variable to convert (`ts`) and the string that should hold the result of the operation (`output`). It converts `ts` to its textual representation according to the SQL standard, which is be `YYYY-MM-DD HH:MM:SS`.

Upon success, the function returns 0 and a negative value if an error occurred.

dttomfmtasc

Convert a timestamp variable to a C char* using a format mask.

```c
int dttomfmtasc(timestamp *ts, char *output, int str_len, char *fmtstr);
```

The function receives a pointer to the timestamp to convert as its first argument (`ts`), a pointer to the output buffer (`output`), the maximal length that has been allocated for the output buffer (`str_len`) and the format mask to use for the conversion (`fmtstr`).

Upon success, the function returns 0 and a negative value if an error occurred.

Internally, this function uses the `PGTYPEStimestamp_fmt_asc` function. See the reference there for information on what format mask specifiers can be used.

intoasc

Convert an interval variable to a C char* string.

```c
int intoasc(interval *i, char *str);
```

The function receives a pointer to the interval variable to convert (`i`) and the string that should hold the result of the operation (`str`). It converts `i` to its textual representation according to the SQL standard, which is be `YYYY-MM-DD HH:MM:SS`.

Upon success, the function returns 0 and a negative value if an error occurred.

rfmtlong

Convert a long integer value to its textual representation using a format mask.

```c
int rfmtlong(long lng_val, char *fmt, char *outbuf);
```

The function receives the long value `lng_val`, the format mask `fmt` and a pointer to the output buffer `outbuf`. It converts the long value according to the format mask to its textual representation.

The format mask can be composed of the following format specifying characters:

- * (asterisk) - if this position would be blank otherwise, fill it with an asterisk.
• & (ampersand) - if this position would be blank otherwise, fill it with a zero.

• # - turn leading zeroes into blanks.

• < - left-justify the number in the string.

• , (comma) - group numbers of four or more digits into groups of three digits separated by a comma.

• . (period) - this character separates the whole-number part of the number from the fractional part.

• - (minus) - the minus sign appears if the number is a negative value.

• + (plus) - the plus sign appears if the number is a positive value.

• ( - this replaces the minus sign in front of the negative number. The minus sign will not appear.

• ) - this character replaces the minus and is printed behind the negative value.

• $ - the currency symbol.

rupshift
Convert a string to upper case.

```c
void rupshift(char *str);
```

The function receives a pointer to the string and transforms every lower case character to upper case.

byleng
Return the number of characters in a string without counting trailing blanks.

```c
int byleng(char *str, int len);
```

The function expects a fixed-length string as its first argument (`str`) and its length as its second argument (`len`). It returns the number of significant characters, that is the length of the string without trailing blanks.

ldchar
Copy a fixed-length string into a null-terminated string.

```c
void ldchar(char *src, int len, char *dest);
```

The function receives the fixed-length string to copy (`src`), its length (`len`) and a pointer to the destination memory (`dest`). Note that you need to reserve at least `len+1` bytes for the string that `dest` points to. The function copies at most `len` bytes to the new location (less if the source string has trailing blanks) and adds the null-terminator.

rgetmsg

```c
int rgetmsg(int msgnum, char *s, int maxsize);
```

This function exists but is not implemented at the moment!
rtyalign

int rtyalign(int offset, int type);

This function exists but is not implemented at the moment!

rtypsize

int rtypsize(int type, int len);

This function exists but is not implemented at the moment!

rtypwidth

int rtpwidth(int sqltype, int sqllen);

This function exists but is not implemented at the moment!

rsetnull

Set a variable to NULL.

int rsetnull(int t, char *ptr);

The function receives an integer that indicates the type of the variable and a pointer to the variable itself that is cast to a C char* pointer.

The following types exist:
• CCHARTYPE - For a variable of type char or char*
• CSHORTTYPE - For a variable of type short int
• CINTTYPE - For a variable of type int
• CBOOLTYPE - For a variable of type boolean
• CFLOATTYPE - For a variable of type float
• CLONGTYPE - For a variable of type long
• CDOUBLETYPE - For a variable of type double
• CDECIMALTYPE - For a variable of type decimal
• CDATETYPE - For a variable of type date
• CDTIMETYPE - For a variable of type timestamp

Here is an example of a call to this function:

$char c[] = "abc       ";
$short s = 17;
$int i = -74874;

rsetnull(CCHARTYPE, (char *) c);
rsetnull(CSHORTTYPE, (char *) &s);
rsetnull(CINTTYPE, (char *) &i);
risnull

Test if a variable is NULL.

int risnull(int t, char *ptr);

The function receives the type of the variable to test (t) as well a pointer to this variable (ptr). Note that the latter needs to be cast to a char*. See the function rsetnull for a list of possible variable types.

Here is an example of how to use this function:

```
char c[] = "abc       ";
short s = 17;
int i = -74874;

risnull(CCHARTYPE, (char *) c);
risnull(CSHORTTYPE, (char *) &s);
risnull(CINTTYPE, (char *) &i);
```

### 36.15.5. Additional Constants

Note that all constants here describe errors and all of them are defined to represent negative values. In the descriptions of the different constants you can also find the value that the constants represent in the current implementation. However you should not rely on this number. You can however rely on the fact all of them are defined to represent negative values.

**ECPG_INFORMIX_NUM_OVERFLOW**

Functions return this value if an overflow occurred in a calculation. Internally it is defined as -1200 (the Informix definition).

**ECPG_INFORMIX_NUM_UNDERFLOW**

Functions return this value if an underflow occurred in a calculation. Internally it is defined as -1201 (the Informix definition).

**ECPG_INFORMIX_DIVIDE_ZERO**

Functions return this value if an attempt to divide by zero is observed. Internally it is defined as -1202 (the Informix definition).

**ECPG_INFORMIX_BAD_YEAR**

Functions return this value if a bad value for a year was found while parsing a date. Internally it is defined as -1204 (the Informix definition).

**ECPG_INFORMIX_BAD_MONTH**

Functions return this value if a bad value for a month was found while parsing a date. Internally it is defined as -1205 (the Informix definition).

**ECPG_INFORMIX_BAD_DAY**

Functions return this value if a bad value for a day was found while parsing a date. Internally it is defined as -1206 (the Informix definition).
ECPG_INFORMIX_ENOSHORTDATE

Functions return this value if a parsing routine needs a short date representation but did not get the date string in the right length. Internally it is defined as -1209 (the Informix definition).

ECPG_INFORMIX_DATE_CONVERT

Functions return this value if an error occurred during date formatting. Internally it is defined as -1210 (the Informix definition).

ECPG_INFORMIX_OUT_OF_MEMORY

Functions return this value if memory was exhausted during their operation. Internally it is defined as -1211 (the Informix definition).

ECPG_INFORMIX_ENOTDMY

Functions return this value if a parsing routine was supposed to get a format mask (like mmddyy) but not all fields were listed correctly. Internally it is defined as -1212 (the Informix definition).

ECPG_INFORMIX_BAD_NUMERIC

Functions return this value either if a parsing routine cannot parse the textual representation for a numeric value because it contains errors or if a routine cannot complete a calculation involving numeric variables because at least one of the numeric variables is invalid. Internally it is defined as -1213 (the Informix definition).

ECPG_INFORMIX_BAD_EXPONENT

Functions return this value if a parsing routine cannot parse an exponent. Internally it is defined as -1216 (the Informix definition).

ECPG_INFORMIX_BAD_DATE

Functions return this value if a parsing routine cannot parse a date. Internally it is defined as -1218 (the Informix definition).

ECPG_INFORMIX_EXTRA_CHARS

Functions return this value if a parsing routine is passed extra characters it cannot parse. Internally it is defined as -1264 (the Informix definition).

36.16. Internals

This section explains how ECPG works internally. This information can occasionally be useful to help users understand how to use ECPG.

The first four lines written by `ecpg` to the output are fixed lines. Two are comments and two are include lines necessary to interface to the library. Then the preprocessor reads through the file and writes output. Normally it just echoes everything to the output.

When it sees an `EXEC SQL` statement, it intervenes and changes it. The command starts with `EXEC SQL` and ends with `;`. Everything in between is treated as an SQL statement and parsed for variable substitution.

Variable substitution occurs when a symbol starts with a colon (:). The variable with that name is looked up among the variables that were previously declared within a `EXEC SQL DECLARE` section.

The most important function in the library is `ECPGdo`, which takes care of executing most commands. It takes a variable number of arguments. This can easily add up to 50 or so arguments, and we hope this will not be a problem on any platform.
The arguments are:

A line number

This is the line number of the original line; used in error messages only.

A string

This is the SQL command that is to be issued. It is modified by the input variables, i.e., the variables that where not known at compile time but are to be entered in the command. Where the variables should go the string contains ?.

Input variables

Every input variable causes ten arguments to be created. (See below.)

\textit{ECPGt\_EOIT}

An \texttt{enum} telling that there are no more input variables.

Output variables

Every output variable causes ten arguments to be created. (See below.) These variables are filled by the function.

\textit{ECPGt\_EORT}

An \texttt{enum} telling that there are no more variables.

For every variable that is part of the SQL command, the function gets ten arguments:

1. The type as a special symbol.
2. A pointer to the value or a pointer to the pointer.
3. The size of the variable if it is a \texttt{char} or \texttt{varchar}.
4. The number of elements in the array (for array fetches).
5. The offset to the next element in the array (for array fetches).
6. The type of the indicator variable as a special symbol.
7. A pointer to the indicator variable.
8. 0
9. The number of elements in the indicator array (for array fetches).
10. The offset to the next element in the indicator array (for array fetches).

Note that not all SQL commands are treated in this way. For instance, an open cursor statement like:

\begin{verbatim}
EXEC SQL OPEN cursor;
\end{verbatim}

is not copied to the output. Instead, the cursor's \texttt{DECLARE} command is used at the position of the \texttt{OPEN} command because it indeed opens the cursor.

Here is a complete example describing the output of the preprocessor of a file \texttt{foo.pgc} (details might change with each particular version of the preprocessor):
EXEC SQL BEGIN DECLARE SECTION;
int index;
int result;
EXEC SQL END DECLARE SECTION;
...
EXEC SQL SELECT res INTO :result FROM mytable WHERE index = :index;

is translated into:

/* Processed by ecpg (2.6.0) */
/* These two include files are added by the preprocessor */
#include <ecpgtype.h>;
#include <ecpglib.h>;

/* exec sql begin declare section */

#line 1 "foo.pgc"

int index;
int result;
/* exec sql end declare section */
...
ECPGdo(__LINE__, NULL, "SELECT res FROM mytable WHERE index = ? ",
ECPT_int,&(index),1L,1L,sizeof(int),
ECPT_NO_INDICATOR, NULL , 0L, 0L, ECPT_EOIT,
ECPT_int,&(result),1L,1L,sizeof(int),
ECPT_NO_INDICATOR, NULL , 0L, 0L, ECPT_EORT);
#line 147 "foo.pgc"

(The indentation here is added for readability and not something the preprocessor does.)
Chapter 37. The Information Schema

The information schema consists of a set of views that contain information about the objects defined in the current database. The information schema is defined in the SQL standard and can therefore be expected to be portable and remain stable — unlike the system catalogs, which are specific to PostgreSQL and are modeled after implementation concerns. The information schema views do not, however, contain information about PostgreSQL-specific features; to inquire about those you need to query the system catalogs or other PostgreSQL-specific views.

Note

When querying the database for constraint information, it is possible for a standard-compliant query that expects to return one row to return several. This is because the SQL standard requires constraint names to be unique within a schema, but PostgreSQL does not enforce this restriction. PostgreSQL automatically-generated constraint names avoid duplicates in the same schema, but users can specify such duplicate names.

This problem can appear when querying information schema views such as check_constraint_routine_usage, check_constraints, domain_constraints, and referential_constraints. Some other views have similar issues but contain the table name to help distinguish duplicate rows, e.g., constraint_column_usage, constraint_table_usage, table_constraints.

37.1. The Schema

The information schema itself is a schema named information_schema. This schema automatically exists in all databases. The owner of this schema is the initial database user in the cluster, and that user naturally has all the privileges on this schema, including the ability to drop it (but the space savings achieved by that are minuscule).

By default, the information schema is not in the schema search path, so you need to access all objects in it through qualified names. Since the names of some of the objects in the information schema are generic names that might occur in user applications, you should be careful if you want to put the information schema in the path.

37.2. Data Types

The columns of the information schema views use special data types that are defined in the information schema. These are defined as simple domains over ordinary built-in types. You should not use these types for work outside the information schema, but your applications must be prepared for them if they select from the information schema.

These types are:

- **cardinal_number**
  
  A nonnegative integer.

- **character_data**
  
  A character string (without specific maximum length).
The Information Schema

sql_identifier

A character string. This type is used for SQL identifiers, the type character_data is used for any other kind of text data.

time_stamp

A domain over the type timestamp with time zone

yes_or_no

A character string domain that contains either YES or NO. This is used to represent Boolean (true/false) data in the information schema. (The information schema was invented before the type boolean was added to the SQL standard, so this convention is necessary to keep the information schema backward compatible.)

Every column in the information schema has one of these five types.

37.3. information_schema_catalog_name

information_schema_catalog_name is a table that always contains one row and one column containing the name of the current database (current catalog, in SQL terminology).

Table 37.1. information_schema_catalog_name Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>catalog_name</td>
<td>sql_identifier</td>
<td>Name of the database that contains this information schema</td>
</tr>
</tbody>
</table>

37.4. administrable_role_authorizations

The view administrable_role_authorizations identifies all roles that the current user has the admin option for.

Table 37.2. administrable_role_authorizations Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role to which this role membership was granted (can be the current user, or a different role in case of nested role memberships)</td>
</tr>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>Always YES</td>
</tr>
</tbody>
</table>

37.5. applicable_roles

The view applicable_roles identifies all roles whose privileges the current user can use. This means there is some chain of role grants from the current user to the role in question. The current user itself is also an applicable role. The set of applicable roles is generally used for permission checking.
Table 37.3. applicable_roles Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role to which this role membership was granted (can be the current user, or a different role in case of nested role memberships)</td>
</tr>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the grantee has the admin option on the role, NO if not</td>
</tr>
</tbody>
</table>

37.6. attributes

The view attributes contains information about the attributes of composite data types defined in the database. (Note that the view does not give information about table columns, which are sometimes called attributes in PostgreSQL contexts.) Only those attributes are shown that the current user has access to (by way of being the owner of or having some privilege on the type).

Table 37.4. attributes Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the data type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the data type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type</td>
</tr>
<tr>
<td>attribute_name</td>
<td>sql_identifier</td>
<td>Name of the attribute</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the attribute within the data type (count starts at 1)</td>
</tr>
<tr>
<td>attribute_default</td>
<td>character_data</td>
<td>Default expression of the attribute</td>
</tr>
<tr>
<td>is_nullable</td>
<td>yes_or_no</td>
<td>YES if the attribute is possibly nullable, NO if it is known not nullable.</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the attribute, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in attribute_udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character type, the maximum possible length in octets (bytes) of</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------</td>
<td>--------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the attribute (always the current database), null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the attribute, null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the attribute, null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a numeric type, this column contains the (declared or implicit) precision of the type for this attribute. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If data_type identifies a numeric type, this column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If data_type identifies an exact numeric type, this column contains the (declared or implicit) scale of the type for this attribute. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------</td>
<td>----------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this attribute, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>If <code>data_type</code> identifies an interval type, this column contains the specification which fields the intervals include for this attribute, e.g., YEAR TO MONTH, DAY TO SECOND, etc. If no field restrictions were specified (that is, the interval accepts all fields), and for all other data types, this field is null.</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see <code>datetime_precision</code> for the fractional seconds precision of interval type attributes)</td>
</tr>
<tr>
<td>attribute_udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the attribute data type is defined in (always the current database)</td>
</tr>
<tr>
<td>attribute_udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the attribute data type is defined in</td>
</tr>
<tr>
<td>attribute_udt_name</td>
<td>sql_identifier</td>
<td>Name of the attribute data type</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the column, unique among the data type descriptors pertaining to the table. This is mainly useful for joining with other instances of such identifiers. (The specific format of the</td>
</tr>
</tbody>
</table>
The view `character_sets` identifies the character sets available in the current database. Since PostgreSQL does not support multiple character sets within one database, this view only shows one, which is the database encoding.

Take note of how the following terms are used in the SQL standard:

character repertoire

An abstract collection of characters, for example UNICODE, UCS, or LATIN1. Not exposed as an SQL object, but visible in this view.

character encoding form

An encoding of some character repertoire. Most older character repertoires only use one encoding form, and so there are no separate names for them (e.g., LATIN1 is an encoding form applicable to the LATIN1 repertoire). But for example Unicode has the encoding forms UTF8, UTF16, etc. (not all supported by PostgreSQL). Encoding forms are not exposed as an SQL object, but are visible in this view.

character set

A named SQL object that identifies a character repertoire, a character encoding, and a default collation. A predefined character set would typically have the same name as an encoding form, but users could define other names. For example, the character set UTF8 would typically identify the character repertoire UCS, encoding form UTF8, and some default collation.

You can think of an “encoding” in PostgreSQL either as a character set or a character encoding form. They will have the same name, and there can only be one in one database.

### Table 37.5. `character_sets` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null.</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null.</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Name of the character set, currently implemented as showing the name of the database encoding.</td>
</tr>
<tr>
<td>character_repertoire</td>
<td>sql_identifier</td>
<td>Character repertoire, showing UCS if the encoding is UTF8, else just the encoding name.</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>form_of_use</td>
<td>sql_identifier</td>
<td>Character encoding form, same as the database encoding</td>
</tr>
<tr>
<td>default_collate_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the default collation (always the current database, if any collation is identified)</td>
</tr>
<tr>
<td>default_collate_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the default collation</td>
</tr>
<tr>
<td>default_collate_name</td>
<td>sql_identifier</td>
<td>Name of the default collation. The default collation is identified as the collation that matches the <code>COLLATE</code> and <code>CTYPE</code> settings of the current database. If there is no such collation, then this column and the associated schema and catalog columns are null.</td>
</tr>
</tbody>
</table>

#### 37.8. check_constraint_routine_usage

The view `check_constraint_routine_usage` identifies routines (functions and procedures) that are used by a check constraint. Only those routines are shown that are owned by a currently enabled role.

**Table 37.6. check_constraint_routine_usage Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
</tbody>
</table>

#### 37.9. check_constraints

The view `check_constraints` contains all check constraints, either defined on a table or on a domain, that are owned by a currently enabled role. (The owner of the table or domain is the owner of the constraint.)
Table 37.7. check_constraints Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>check_clause</td>
<td>character_data</td>
<td>The check expression of the check constraint</td>
</tr>
</tbody>
</table>

37.10. collations

The view collations contains the collations available in the current database.

Table 37.8. collations Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation (always the current database)</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the default collation</td>
</tr>
<tr>
<td>pad_attribute</td>
<td>character_data</td>
<td>Always NO PAD (The alternative PAD SPACE is not supported by PostgreSQL.)</td>
</tr>
</tbody>
</table>

37.11. collation_character_set_applicability

The view collation_character_set_applicability identifies which character set the available collations are applicable to. In PostgreSQL, there is only one character set per database (see explanation in Section 37.7), so this view does not provide much useful information.

Table 37.9. collation_character_set_applicability Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation (always the current database)</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the default collation</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null</td>
</tr>
</tbody>
</table>
The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Name of the character set</td>
</tr>
</tbody>
</table>

### 37.12. column_domain_usage

The view `column_domain_usage` identifies all columns (of a table or a view) that make use of some domain defined in the current database and owned by a currently enabled role.

**Table 37.10. column_domain_usage Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
</tbody>
</table>

### 37.13. column_options

The view `column_options` contains all the options defined for foreign table columns in the current database. Only those foreign table columns are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.11. column_options Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

### 37.14. column_privileges

The view `column_privileges` identifies all privileges granted on columns to a currently enabled role or by a currently enabled role. There is one row for each combination of column, grantor, and grantee.
If a privilege has been granted on an entire table, it will show up in this view as a grant for each column, but only for the privilege types where column granularity is possible: SELECT, INSERT, UPDATE, REFERENCES.

### Table 37.12. column_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, or REFERENCES</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

### 37.15. column_udt_usage

The view column_udt_usage identifies all columns that use data types owned by a currently enabled role. Note that in PostgreSQL, built-in data types behave like user-defined types, so they are included here as well. See also Section 37.16 for details.

### Table 37.13. column_udt_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the column data type (the underlying type of the domain, if applicable) is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the column data type (the underlying type of the domain, if applicable) is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the column data type (the underlying type of the domain, if applicable)</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
</tbody>
</table>
The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
</tbody>
</table>

### 37.16. columns

The view `columns` contains information about all table columns (or view columns) in the database. System columns (oid, etc.) are not included. Only those columns are shown that the current user has access to (by way of being the owner or having some privilege).

#### Table 37.14. columns Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the column within the table (count starts at 1)</td>
</tr>
<tr>
<td>column_default</td>
<td>character_data</td>
<td>Default expression of the column</td>
</tr>
<tr>
<td>is_nullable</td>
<td>yes_or_no</td>
<td>YES if the column is possibly nullable, NO if it is known not nullable. A not-null constraint is one way a column can be known not nullable, but there can be others.</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the column, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in udt_name and associated columns). If the column is based on a domain, this column refers to the type underlying the domain (and the domain is identified in domain_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character type, the maximum possible length in octets (bytes) of a datum; null for all other data types. The maximum octet...</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------</td>
<td>--------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>length</td>
<td>length</td>
<td>length depends on the declared character maximum length (see above) and the server encoding.</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a numeric type, this column contains the (declared or implicit) precision of the type for this column. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If data_type identifies a numeric type, this column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If data_type identifies an exact numeric type, this column contains the (declared or implicit) scale of the type for this column. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this column, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
</tbody>
</table>
| interval_type           | character_data     | If data_type identifies an interval type, this column contains the specification which fields the intervals include for this column, e.g., YEAR TO MONTH, DAY TO SECOND, etc. If no field restrictions were specified (that
<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see date-time_precision for the fractional seconds precision of interval type columns)</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the column (always the current database), null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the column, null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the column, null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the database that the domain is defined in (always the current database), else null.</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the schema that the domain is defined in, else null.</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the domain, else null.</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the column data type (the underlying type of the domain, if applicable) is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the column data type (the underlying type of the domain, if applicable) is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the column data type (the underlying type of the domain, if applicable)</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------</td>
<td>-------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the column, unique among the data type descriptors pertaining to the table. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>is_self_referencing</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_identity</td>
<td>yes_or_no</td>
<td>If the column is an identity column, then YES, else NO.</td>
</tr>
<tr>
<td>identity_generation</td>
<td>character_data</td>
<td>If the column is an identity column, then ALWAYS or BY DEFAULT, reflecting the definition of the column.</td>
</tr>
<tr>
<td>identity_start</td>
<td>character_data</td>
<td>If the column is an identity column, then the start value of the internal sequence, else null.</td>
</tr>
<tr>
<td>identity_increment</td>
<td>character_data</td>
<td>If the column is an identity column, then the increment of the internal sequence, else null.</td>
</tr>
<tr>
<td>identity_maximum</td>
<td>character_data</td>
<td>If the column is an identity column, then the maximum value of the internal sequence, else null.</td>
</tr>
<tr>
<td>identity_minimum</td>
<td>character_data</td>
<td>If the column is an identity column, then the minimum value of the internal sequence, else null.</td>
</tr>
<tr>
<td>identity_cycle</td>
<td>yes_or_no</td>
<td>If the column is an identity column, then YES if the internal sequence cycles or NO if it does not; otherwise null.</td>
</tr>
<tr>
<td>is_generated</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>generation_expression</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_updatable</td>
<td>yes_or_no</td>
<td>YES if the column is updatable, NO if not (Columns in base tables are always updatable, columns in views not necessarily)</td>
</tr>
</tbody>
</table>
Since data types can be defined in a variety of ways in SQL, and PostgreSQL contains additional ways to define data types, their representation in the information schema can be somewhat difficult. The column `data_type` is supposed to identify the underlying built-in type of the column. In PostgreSQL, this means that the type is defined in the system catalog schema `pg_catalog`. This column might be useful if the application can handle the well-known built-in types specially (for example, format the numeric types differently or use the data in the precision columns). The columns `udt_name`, `udt_schema`, and `udt_catalog` always identify the underlying data type of the column, even if the column is based on a domain. (Since PostgreSQL treats built-in types like user-defined types, built-in types appear here as well. This is an extension of the SQL standard.) These columns should be used if an application wants to process data differently according to the type, because in that case it wouldn’t matter if the column is really based on a domain. If the column is based on a domain, the identity of the domain is stored in the columns `domain_name`, `domain_schema`, and `domain_catalog`. If you want to pair up columns with their associated data types and treat domains as separate types, you could write `coalesce(domain_name, udt_name)`, etc.

### 37.17. `constraint_column_usage`

The view `constraint_column_usage` identifies all columns in the current database that are used by some constraint. Only those columns are shown that are contained in a table owned by a currently enabled role. For a check constraint, this view identifies the columns that are used in the check expression. For a foreign key constraint, this view identifies the columns that the foreign key references. For a unique or primary key constraint, this view identifies the constrained columns.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>table_catalog</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the table that contains the column that is used by some constraint (always the current database)</td>
</tr>
<tr>
<td><code>table_schema</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the schema that contains the table that contains the column that is used by some constraint</td>
</tr>
<tr>
<td><code>table_name</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the table that contains the column that is used by some constraint</td>
</tr>
<tr>
<td><code>column_name</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the column that is used by some constraint</td>
</tr>
<tr>
<td><code>constraint_catalog</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td><code>constraint_schema</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td><code>constraint_name</code></td>
<td><code>sql_identifier</code></td>
<td>Name of the constraint</td>
</tr>
</tbody>
</table>

### 37.18. `constraint_table_usage`

The view `constraint_table_usage` identifies all tables in the current database that are used by some constraint and are owned by a currently enabled role. (This is different from the view `table_constraints`, which identifies all table constraints along with the table they are defined on.) For a foreign key constraint, this view identifies the table that the foreign key references. For a unique
or primary key constraint, this view simply identifies the table the constraint belongs to. Check constraints and not-null constraints are not included in this view.

Table 37.16. constraint_table_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that is used by some constraint (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that is used by some constraint</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that is used by some constraint</td>
</tr>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
</tbody>
</table>

37.19. data_type_privileges

The view data_type_privileges identifies all data type descriptors that the current user has access to, by way of being the owner of the described object or having some privilege for it. A data type descriptor is generated whenever a data type is used in the definition of a table column, a domain, or a function (as parameter or return type) and stores some information about how the data type is used in that instance (for example, the declared maximum length, if applicable). Each data type descriptor is assigned an arbitrary identifier that is unique among the data type descriptor identifiers assigned for one object (table, domain, function). This view is probably not useful for applications, but it is used to define some other views in the information schema.

Table 37.17. data_type_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the described object (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the described object</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the described object</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>The type of the described object: one of TABLE (the data type descriptor pertains to a column of that table), DOMAIN (the data type descriptors pertains to that domain), ROUTINE (the data type descriptor pertains to a parameter or the return data type of that function).</td>
</tr>
</tbody>
</table>
| dtd_identifier    | sql_identifier| The identifier of the data type descriptor, which is unique among
The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>the data type descriptors for that same object.</td>
</tr>
</tbody>
</table>

### 37.20. domain_constraints

The view `domain_constraints` contains all constraints belonging to domains defined in the current database. Only those domains are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.18. domain_constraints Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
<tr>
<td>is_deferrable</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable, NO if not</td>
</tr>
<tr>
<td>initially_deferred</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable and initially deferred, NO if not</td>
</tr>
</tbody>
</table>

### 37.21. domain_udt_usage

The view `domain_udt_usage` identifies all domains that are based on data types owned by a currently enabled role. Note that in PostgreSQL, built-in data types behave like user-defined types, so they are included here as well.

**Table 37.19. domain_udt_usage Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the domain data type is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the domain data type is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the domain data type</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Name of the schema that contains the domain</td>
</tr>
</tbody>
</table>

985
The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
</tbody>
</table>

### 37.22. domains

The view `domains` contains all domains defined in the current database. Only those domains are shown that the current user has access to (by way of being the owner or having some privilege).

#### Table 37.20. domains Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the domain, if it is a built-in type, or ARRAY if it is some array (in that case, see the view <code>element_types</code>), else USER-DEFINED (in that case, the type is identified in <code>udt_name</code> and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If the domain has a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If the domain has a character type, the maximum possible length in octets (bytes) of a datum; null for all other data types. The maximum octet length depends on the declared character maximum length (see above) and the server encoding.</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the domain (always the current database), null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the domain, null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>----------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the domain, null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If the domain has a numeric type, this column contains the (declared or implicit) precision of the type for this domain. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If the domain has a numeric type, this column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If the domain has an exact numeric type, this column contains the (declared or implicit) scale of the type for this domain. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this domain, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>If data_type identifies an interval type, this column contains the specification which fields the intervals include for this domain, e.g., YEAR TO MONTH, DAY TO SECOND, etc. If no field restrictions were specified (that</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>-----------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see date-time_precision for the fractional seconds precision of interval type domains)</td>
</tr>
<tr>
<td>domain_default</td>
<td>character_data</td>
<td>Default expression of the domain</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the domain data type is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the domain data type is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the domain data type</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the domain, unique among the data type descriptors pertaining to the domain (which is trivial, because a domain only contains one data type descriptor). This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
</tbody>
</table>

### 37.23. `element_types`

The view `element_types` contains the data type descriptors of the elements of arrays. When a table column, composite-type attribute, domain, function parameter, or function return value is defined to be of an array type, the respective information schema view only contains `ARRAY` in the column `data_type`. To obtain information on the element type of the array, you can join the respective view with this view. For example, to show the columns of a table with data types and array element types, if applicable, you could do:

```sql
SELECT c.column_name, c.data_type, e.data_type AS element_type
FROM information_schema.columns c LEFT JOIN
     information_schema.element_types e
```
ON ((c.table_catalog, c.table_schema, c.table_name, 'TABLE', c.dtd_identifier) = (e.object_catalog, e.object_schema, e.object_name, e.object_type, e.collection_type_identifier))
WHERE c.table_schema = '...' AND c.table_name = '...'
ORDER BY c.ordinal_position;

This view only includes objects that the current user has access to, by way of being the owner or having some privilege.

### Table 37.21. *element_types* Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the object that uses the array being described (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the object that uses the array being described</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the object that uses the array being described</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>The type of the object that uses the array being described: one of TABLE (the array is used by a column of that table), USER-DEFINED TYPE (the array is used by an attribute of that composite type), DOMAIN (the array is used by that domain), ROUTINE (the array is used by a parameter or the return data type of that function).</td>
</tr>
<tr>
<td>collection_type_identifier</td>
<td>sql_identifier</td>
<td>The identifier of the data type descriptor of the array being described. Use this to join with the dtd_identifier columns of other information schema views.</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the array elements, if it is a built-in type, else USER-DEFINED (in that case, the type is identified in udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>-----------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the element type (always the current database), null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the element type, null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the element type, null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>domain_default</td>
<td>character_data</td>
<td>Not yet implemented</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the data type of the elements is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the data type of the elements is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type of the elements</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the element. This is currently not useful.</td>
</tr>
</tbody>
</table>

#### 37.24. enabled_roles

The view `enabled_roles` identifies the currently “enabled roles”. The enabled roles are recursively defined as the current user together with all roles that have been granted to the enabled roles with automatic inheritance. In other words, these are all roles that the current user has direct or indirect, automatically inheriting membership in.

For permission checking, the set of “applicable roles” is applied, which can be broader than the set of enabled roles. So generally, it is better to use the view `applicable_roles` instead of this one; See Section 37.5 for details on `applicable_roles` view.

**Table 37.22. enabled_roles Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
</tbody>
</table>

#### 37.25. foreign_data_wrapper_options

The view `foreign_data_wrapper_options` contains all the options defined for foreign-data wrappers in the current database. Only those foreign-data wrappers are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.23. foreign_data_wrapper_options Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_data_wrapper_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign-data wrapper is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_data_wrapper_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

#### 37.26. foreign_data_wrappers

The view `foreign_data_wrappers` contains all foreign-data wrappers defined in the current database. Only those foreign-data wrappers are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.24. foreign_data_wrappers Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_data_wrapper_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign-data wrapper (always the current database)</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_data_wrap-per_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the owner of the foreign server</td>
</tr>
<tr>
<td>library_name</td>
<td>character_data</td>
<td>File name of the library that implementing this foreign-data wrapper</td>
</tr>
<tr>
<td>foreign_data_wrap-per_language</td>
<td>character_data</td>
<td>Language used to implement this foreign-data wrapper</td>
</tr>
</tbody>
</table>

#### 37.27. foreign_server_options

The view `foreign_server_options` contains all the options defined for foreign servers in the current database. Only those foreign servers are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.25. foreign_server_options Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

#### 37.28. foreign_servers

The view `foreign_servers` contains all foreign servers defined in the current database. Only those foreign servers are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.26. foreign_servers Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>foreign_data_wrap-per_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign-data wrapper used by the foreign server (always the current database)</td>
</tr>
<tr>
<td>foreign_data_wrap-per_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper used by the foreign server</td>
</tr>
<tr>
<td>foreign_server_type</td>
<td>character_data</td>
<td>Foreign server type information, if specified upon creation</td>
</tr>
<tr>
<td>foreign_server_version</td>
<td>character_data</td>
<td>Foreign server version information, if specified upon creation</td>
</tr>
</tbody>
</table>

992
### 37.29. `foreign_table_options`

The view `foreign_table_options` contains all the options defined for foreign tables in the current database. Only those foreign tables are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.27. `foreign_table_options` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign table (always the current database)</td>
</tr>
<tr>
<td>foreign_table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
<tr>
<td>foreign_table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

### 37.30. `foreign_tables`

The view `foreign_tables` contains all foreign tables defined in the current database. Only those foreign tables are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 37.28. `foreign_tables` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign table is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
<tr>
<td>foreign_table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
</tbody>
</table>

### 37.31. `key_column_usage`

The view `key_column_usage` identifies all columns in the current database that are restricted by some unique, primary key, or foreign key constraint. Check constraints are not included in this view. Only those columns are shown that the current user has access to, by way of being the owner or having some privilege.
Table 37.29. **key_column_usage** Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the column that is restricted by this constraint (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column that is restricted by this constraint</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column that is restricted by this constraint</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column that is restricted by this constraint</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the column within the constraint key (count starts at 1)</td>
</tr>
<tr>
<td>position_in_unique_constraint</td>
<td>cardinal_number</td>
<td>For a foreign-key constraint, ordinal position of the referenced column within its unique constraint (count starts at 1); otherwise null</td>
</tr>
</tbody>
</table>

37.32. **parameters**

The view **parameters** contains information about the parameters (arguments) of all functions in the current database. Only those functions are shown that the current user has access to (by way of being the owner or having some privilege).

Table 37.30. **parameters** Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the parameter in the argument list of the function (count starts at 1)</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>--------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>parameter_mode</td>
<td>character_data</td>
<td>IN for input parameter, OUT for output parameter, and INOUT for input/output parameter.</td>
</tr>
<tr>
<td>is_result</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>parameter_name</td>
<td>sql_identifier</td>
<td>Name of the parameter, or null if the parameter has no name.</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the parameter, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL.</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the data type of the parameter is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the data type of the parameter is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type of the parameter</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL.</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the parameter, unique among the data type descriptors pertaining to the function. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>parameter_default</td>
<td>character_data</td>
<td>The default expression of the parameter, or null if none or if the function is not owned by a currently enabled role.</td>
</tr>
</tbody>
</table>

### 37.33. referential_constraints

The view `referential_constraints` contains all referential (foreign key) constraints in the current database. Only those constraints are shown for which the current user has write access to the referencing table (by way of being the owner or having some privilege other than SELECT).

**Table 37.31. referential_constraints Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>unique_constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the unique or primary key constraint</td>
</tr>
<tr>
<td></td>
<td></td>
<td>that the foreign key constraint references (always the current database)</td>
</tr>
<tr>
<td>unique_constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the unique or primary key constraint</td>
</tr>
<tr>
<td></td>
<td></td>
<td>that the foreign key constraint references</td>
</tr>
<tr>
<td>unique_constraint_name</td>
<td>sql_identifier</td>
<td>Name of the unique or primary key constraint that the foreign key constraint</td>
</tr>
<tr>
<td></td>
<td></td>
<td>references</td>
</tr>
<tr>
<td>match_option</td>
<td>character_data</td>
<td>Match option of the foreign key constraint: FULL, PARTIAL, or NONE.</td>
</tr>
<tr>
<td>update_rule</td>
<td>character_data</td>
<td>Update rule of the foreign key constraint: CASCADE, SET NULL, SET DEFAULT,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>RESTRICT, or NO ACTION.</td>
</tr>
<tr>
<td>delete_rule</td>
<td>character_data</td>
<td>Delete rule of the foreign key constraint: CASCADE, SET NULL, SET DEFAULT,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>RESTRICT, or NO ACTION.</td>
</tr>
</tbody>
</table>

**37.34. role_column_grants**

The view `role_column_grants` identifies all privileges granted on columns where the grantor or grantee is a currently enabled role. Further information can be found under `column_privileges`. The only effective difference between this view and `column_privileges` is that this view omits columns that have been made accessible to the current user by way of a grant to `PUBLIC`.

**Table 37.32. role_column_grants Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
</tbody>
</table>
The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, or REFERENCES</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

37.35. role_routine_grants

The view role_routine_grants identifies all privileges granted on functions where the grantor or grantee is a currently enabled role. Further information can be found under routine_privileges. The only effective difference between this view and routine_privileges is that this view omits functions that have been made accessible to the current user by way of a grant to PUBLIC.

Table 37.33. role_routine_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
<tr>
<td>routine_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>routine_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>routine_name</td>
<td>sql_identifier</td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always EXECUTE (the only privilege type for functions)</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

37.36. role_table_grants

The view role_table_grants identifies all privileges granted on tables or views where the grantor or grantee is a currently enabled role. Further information can be found under table_privileges. The only effective difference between this view and table_privileges is that this view omits tables that have been made accessible to the current user by way of a grant to PUBLIC.
### Table 37.34. role_table_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, or TRIGGER</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
<tr>
<td>with_hierarchy</td>
<td>yes_or_no</td>
<td>In the SQL standard, WITH HIERARCHY OPTION is a separate (sub-)privilege allowing certain operations on table inheritance hierarchies. In PostgreSQL, this is included in the SELECT privilege, so this column shows YES if the privilege is SELECT, else NO.</td>
</tr>
</tbody>
</table>

### Table 37.35. role_udt_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>The name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>The name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always TYPE USAGE</td>
</tr>
</tbody>
</table>
37.38. role_usage_grants

The view `role_usage_grants` identifies USAGE privileges granted on various kinds of objects where the grantor or grantee is a currently enabled role. Further information can be found under `usage_privileges`. The only effective difference between this view and `usage_privileges` is that this view omits objects that have been made accessible to the current user by way of a grant to PUBLIC.

### Table 37.36. role_usage_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>The name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>The name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the object (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the object, if applicable, else an empty string</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the object</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>COLLATION or DOMAIN or FOREIGN DATA WRAPPER or FOREIGN SERVER or SEQUENCE</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable. NO if not</td>
</tr>
</tbody>
</table>

37.39. routine_privileges

The view `routine_privileges` identifies all privileges granted on functions to a currently enabled role or by a currently enabled role. There is one row for each combination of function, grantor, and grantee.

### Table 37.37. routine_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
<tr>
<td>routine_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>routine_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>routine_name</td>
<td>sql_identifier</td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always EXECUTE (the only privilege type for functions)</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

#### 37.40. routines

The view routines contains all functions and procedures in the current database. Only those functions and procedures are shown that the current user has access to (by way of being the owner or having some privilege).

### Table 37.38. routines Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. This is a name that uniquely identifies the function in the schema, even if the real name of the function is overloaded. The format of the specific name is not defined, it should only be used to compare it to other instances of specific routine names.</td>
</tr>
<tr>
<td>routine_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>routine_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>routine_name</td>
<td>sql_identifier</td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>routine_type</td>
<td>character_data</td>
<td>FUNCTION for a function, PROCEDURE for a procedure</td>
</tr>
<tr>
<td>module_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>--------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>module_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>module_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Return data type of the function, if it is a built-in type, or ARRAY if it is some array (in that case, see the view <code>element_types</code>), else USER-DEFINED (in that case, the type is identified in <code>type_udt_name</code> and associated columns). Null for a procedure.</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>-----------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>type_udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the return data type of the function is defined in (always the current database). Null for a procedure.</td>
</tr>
<tr>
<td>type_udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the return data type of the function is defined in. Null for a procedure.</td>
</tr>
<tr>
<td>type_udt_name</td>
<td>sql_identifier</td>
<td>Name of the return data type of the function. Null for a procedure.</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the return data type of this function, unique among the data type descriptors pertaining to the function. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>routine_body</td>
<td>character_data</td>
<td>If the function is an SQL function, then SQL, else EXTERNAL.</td>
</tr>
<tr>
<td>routine_definition</td>
<td>character_data</td>
<td>The source text of the function (null if the function is not owned by a currently enabled role). (According to the SQL standard, this column is only applicable if routine_body is SQL, but in PostgreSQL it will contain whatever source text was specified when the function was created.)</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------</td>
<td>----------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>external_name</td>
<td>character_data</td>
<td>If this function is a C function, then the external name (link symbol) of the function; else null. (This works out to be the same value that is shown in routine_definition.)</td>
</tr>
<tr>
<td>external_language</td>
<td>character_data</td>
<td>The language the function is written in</td>
</tr>
<tr>
<td>parameter_style</td>
<td>character_data</td>
<td>Always GENERAL (The SQL standard defines other parameter styles, which are not available in PostgreSQL.)</td>
</tr>
<tr>
<td>is_deterministic</td>
<td>yes_or_no</td>
<td>If the function is declared immutable (called deterministic in the SQL standard), then YES, else NO. (You cannot query the other volatility levels available in PostgreSQL through the information schema.)</td>
</tr>
<tr>
<td>sql_data_access</td>
<td>character_data</td>
<td>Always MODIFIES, meaning that the function possibly modifies SQL data. This information is not useful for PostgreSQL.</td>
</tr>
<tr>
<td>is_null_call</td>
<td>yes_or_no</td>
<td>If the function automatically returns null if any of its arguments are null, then YES, else NO. Null for a procedure.</td>
</tr>
<tr>
<td>sql_path</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>schema_level_routine</td>
<td>yes_or_no</td>
<td>Always YES (The opposite would be a method of a user-defined type, which is a feature not available in PostgreSQL.)</td>
</tr>
<tr>
<td>max_dynamic_result_sets</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_user_defined_cast</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>isimplicitly_invocable</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>security_type</td>
<td>character_data</td>
<td>If the function runs with the privileges of the current user, then INVOKER, if the function runs with the privileges of the user who defined it, then DEFINER.</td>
</tr>
<tr>
<td>to_sql_specific_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>to_sql_specific_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>to_sql_specific_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------------------</td>
<td>----------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>created</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>last_altered</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>new_savepoint_level</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_udt_dependent</td>
<td>yes_or_no</td>
<td>Currently always NO. The alternative YES applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>result_cast_from_data_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_max_length</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_octet_length</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collection_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collection_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collection_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_scale</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_datetime_precision</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_interval_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_type_udt_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>
### 37.41. schemata

The view `schemata` contains all schemas in the current database that the current user has access to (by way of being the owner or having some privilege).

#### Table 37.39. schemata Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>catalog_name</td>
<td>sql_identifier</td>
<td>Name of the database that the schema is contained in (always the current database)</td>
</tr>
<tr>
<td>schema_name</td>
<td>sql_identifier</td>
<td>Name of the schema</td>
</tr>
<tr>
<td>schema_owner</td>
<td>sql_identifier</td>
<td>Name of the owner of the schema</td>
</tr>
<tr>
<td>default_character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>default_character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>default_character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>sql_path</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

### 37.42. sequences

The view `sequences` contains all sequences defined in the current database. Only those sequences are shown that the current user has access to (by way of being the owner or having some privilege).

#### Table 37.40. sequences Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sequence_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the sequence (always the current database)</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sequence_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the sequence</td>
</tr>
<tr>
<td>sequence_name</td>
<td>sql_identifier</td>
<td>Name of the sequence</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>The data type of the sequence.</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>This column contains the (declared or implicit) precision of the sequence data type (see above). The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>This column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>This column contains the (declared or implicit) scale of the sequence data type (see above). The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix.</td>
</tr>
<tr>
<td>start_value</td>
<td>character_data</td>
<td>The start value of the sequence</td>
</tr>
<tr>
<td>minimum_value</td>
<td>character_data</td>
<td>The minimum value of the sequence</td>
</tr>
<tr>
<td>maximum_value</td>
<td>character_data</td>
<td>The maximum value of the sequence</td>
</tr>
<tr>
<td>increment</td>
<td>character_data</td>
<td>The increment of the sequence</td>
</tr>
<tr>
<td>cycle_option</td>
<td>yes_or_no</td>
<td>YES if the sequence cycles, else NO</td>
</tr>
</tbody>
</table>

Note that in accordance with the SQL standard, the start, minimum, maximum, and increment values are returned as character strings.

### 37.43. sql_features

The table `sql_features` contains information about which formal features defined in the SQL standard are supported by PostgreSQL. This is the same information that is presented in Appendix D. There you can also find some additional background information.

#### Table 37.41. `sql_features` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>Identifier string of the feature</td>
</tr>
</tbody>
</table>
37.44. sql_imp_plementa_tion_info

The table sql_implementation_info contains information about various aspects that are left implementation-defined by the SQL standard. This information is primarily intended for use in the context of the ODBC interface; users of other interfaces will probably find this information to be of little use. For this reason, the individual implementation information items are not described here; you will find them in the description of the ODBC interface.

Table 37.42. sql_implementation_info Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>implementation_info_id</td>
<td>character_data</td>
<td>Identifier string of the implementation information item</td>
</tr>
<tr>
<td>implementation_info_name</td>
<td>character_data</td>
<td>Descriptive name of the implementation information item</td>
</tr>
<tr>
<td>integer_value</td>
<td>cardinal_number</td>
<td>Value of the implementation information item, or null if the value is contained in the column character_value</td>
</tr>
<tr>
<td>character_value</td>
<td>character_data</td>
<td>Value of the implementation information item, or null if the value is contained in the column integer_value</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment pertaining to the implementation information item</td>
</tr>
</tbody>
</table>

37.45. sql_languages

The table sql_languages contains one row for each SQL language binding that is supported by PostgreSQL. PostgreSQL supports direct SQL and embedded SQL in C; that is all you will learn from this table.

This table was removed from the SQL standard in SQL:2008, so there are no entries referring to standards later than SQL:2003.
Table 37.43. sql_languages Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sql_language_source</td>
<td>character_data</td>
<td>The name of the source of the language definition; always ISO 9075, that is, the SQL standard</td>
</tr>
<tr>
<td>sql_language_year</td>
<td>character_data</td>
<td>The year the standard referenced in sql_language_source was approved.</td>
</tr>
<tr>
<td>sql_language_conformance</td>
<td>character_data</td>
<td>The standard conformance level for the language binding. For ISO 9075:2003 this is always CORE.</td>
</tr>
<tr>
<td>sql_language_integrity</td>
<td>character_data</td>
<td>Always null (This value is relevant to an earlier version of the SQL standard.)</td>
</tr>
<tr>
<td>sql_language_implementation</td>
<td>character_data</td>
<td>Always null</td>
</tr>
<tr>
<td>sql_language_binding_style</td>
<td>character_data</td>
<td>The language binding style, either DIRECT or EMBEDDED</td>
</tr>
<tr>
<td>sql_language_programming_language</td>
<td>character_data</td>
<td>The programming language, if the binding style is EMBEDDED, else null. PostgreSQL only supports the language C.</td>
</tr>
</tbody>
</table>

37.46. sql_packages

The table sql_packages contains information about which feature packages defined in the SQL standard are supported by PostgreSQL. Refer to Appendix D for background information on feature packages.

Table 37.44. sql_packages Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>Identifier string of the package</td>
</tr>
<tr>
<td>feature_name</td>
<td>character_data</td>
<td>Descriptive name of the package</td>
</tr>
<tr>
<td>is_supported</td>
<td>yes_or_no</td>
<td>YES if the package is fully supported by the current version of PostgreSQL, NO if not</td>
</tr>
<tr>
<td>is_verified_by</td>
<td>character_data</td>
<td>Always null, since the PostgreSQL development group does not perform formal testing of feature conformance</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment about the supported status of the package</td>
</tr>
</tbody>
</table>

37.47. sql_parts

The table sql_parts contains information about which of the several parts of the SQL standard are supported by PostgreSQL.
### Table 37.45. sql_parts Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>An identifier string containing the number of the part</td>
</tr>
<tr>
<td>feature_name</td>
<td>character_data</td>
<td>Descriptive name of the part</td>
</tr>
<tr>
<td>is_supported</td>
<td>yes_or_no</td>
<td>YES if the part is fully supported by the current version of PostgreSQL, NO if not</td>
</tr>
<tr>
<td>is_verified_by</td>
<td>character_data</td>
<td>Always null, since the PostgreSQL development group does not perform formal testing of feature conformance</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment about the supported status of the part</td>
</tr>
</tbody>
</table>

### 37.48. sql_sizing

The table `sql_sizing` contains information about various size limits and maximum values in PostgreSQL. This information is primarily intended for use in the context of the ODBC interface; users of other interfaces will probably find this information to be of little use. For this reason, the individual sizing items are not described here; you will find them in the description of the ODBC interface.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sizing_id</td>
<td>cardinal_number</td>
<td>Identifier of the sizing item</td>
</tr>
<tr>
<td>sizing_name</td>
<td>character_data</td>
<td>Descriptive name of the sizing item</td>
</tr>
<tr>
<td>supported_value</td>
<td>cardinal_number</td>
<td>Value of the sizing item, or 0 if the size is unlimited or cannot be determined, or null if the features for which the sizing item is applicable are not supported</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment pertaining to the sizing item</td>
</tr>
</tbody>
</table>

### 37.49. sql_sizing_profiles

The table `sql_sizing_profiles` contains information about the `sql_sizing` values that are required by various profiles of the SQL standard. PostgreSQL does not track any SQL profiles, so this table is empty.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sizing_id</td>
<td>cardinal_number</td>
<td>Identifier of the sizing item</td>
</tr>
<tr>
<td>sizing_name</td>
<td>character_data</td>
<td>Descriptive name of the sizing item</td>
</tr>
<tr>
<td>profile_id</td>
<td>character_data</td>
<td>Identifier string of a profile</td>
</tr>
<tr>
<td>required_value</td>
<td>cardinal_number</td>
<td>The value required by the SQL profile for the sizing item, or 0 if the profile places no limit on the</td>
</tr>
</tbody>
</table>
37.50. **table_constraints**

The view `table_constraints` contains all constraints belonging to tables that the current user owns or has some privilege other than `SELECT` on.

**Table 37.48. **table_constraints** Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td><code>sql_identifier</code></td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td><code>sql_identifier</code></td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>table_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td><code>sql_identifier</code></td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td><code>sql_identifier</code></td>
<td>Name of the table</td>
</tr>
<tr>
<td>constraint_type</td>
<td><code>character_data</code></td>
<td>Type of the constraint: <code>CHECK</code>, <code>FOREIGN KEY</code>, <code>PRIMARY KEY</code>, or <code>UNIQUE</code></td>
</tr>
<tr>
<td>is_deferrable</td>
<td><code>yes_or_no</code></td>
<td>YES if the constraint is deferrable, NO if not</td>
</tr>
<tr>
<td>initially_deferred</td>
<td><code>yes_or_no</code></td>
<td>YES if the constraint is deferrable and initially deferred, NO if not</td>
</tr>
<tr>
<td>enforced</td>
<td><code>yes_or_no</code></td>
<td>Applies to a feature not available in PostgreSQL (currently always YES)</td>
</tr>
</tbody>
</table>

37.51. **table_privileges**

The view `table_privileges` identifies all privileges granted on tables or views to a currently enabled role or by a currently enabled role. There is one row for each combination of table, grantor, and grantee.

**Table 37.49. **table_privileges** Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td><code>sql_identifier</code></td>
<td>Name of the role that granted the privilege</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, or TRIGGER</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
<tr>
<td>with_hierarchy</td>
<td>yes_or_no</td>
<td>In the SQL standard, WITH HIERARCHY OPTION is a separate (sub-)privilege allowing certain operations on table inheritance hierarchies. In PostgreSQL, this is included in the SELECT privilege, so this column shows YES if the privilege is SELECT, else NO.</td>
</tr>
</tbody>
</table>

### 37.52. tables

The view tables contains all tables and views defined in the current database. Only those tables and views are shown that the current user has access to (by way of being the owner or having some privilege).

#### Table 37.50. tables Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>table_type</td>
<td>character_data</td>
<td>Type of the table: BASE TABLE for a persistent base table (the normal table type), VIEW for a view, FOREIGN for a foreign table, or LOCAL TEMPORARY for a temporary table</td>
</tr>
<tr>
<td>self_referencing_column_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>reference_generation</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>user_defined_type_catalog</td>
<td>sql_identifier</td>
<td>If the table is a typed table, the name of the database that contains the column</td>
</tr>
</tbody>
</table>
### The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>user_def_type_schema</td>
<td>sql_identifier</td>
<td>If the table is a typed table, the name of the schema that contains the underlying data type, else null.</td>
</tr>
<tr>
<td>user_def_type_name</td>
<td>sql_identifier</td>
<td>If the table is a typed table, the name of the underlying data type, else null.</td>
</tr>
<tr>
<td>is_insertable_into</td>
<td>yes_or_no</td>
<td>YES if the table is insertable into, NO if not (Base tables are always insertable into, views not necessarily.)</td>
</tr>
<tr>
<td>is_typed</td>
<td>yes_or_no</td>
<td>YES if the table is a typed table, NO if not</td>
</tr>
<tr>
<td>commit_action</td>
<td>character_data</td>
<td>Not yet implemented</td>
</tr>
</tbody>
</table>

### 37.53. transforms

The view transforms contains information about the transforms defined in the current database. More precisely, it contains a row for each function contained in a transform (the “from SQL” or “to SQL” function).

#### Table 37.51. transforms Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the type the transform is for (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the type the transform is for</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type the transform is for</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
<tr>
<td>group_name</td>
<td>sql_identifier</td>
<td>The SQL standard allows defining transforms in “groups”, and selecting a group at run time. PostgreSQL does not support this. Instead, transforms are specific to a language. As a compromise, this field contains the language the transform is for.</td>
</tr>
<tr>
<td>transform_type</td>
<td>character_data</td>
<td>FROM SQL or TO SQL</td>
</tr>
</tbody>
</table>
37.54. triggered_update_columns

For triggers in the current database that specify a column list (like UPDATE OF column1, column2), the view `triggered_update_columns` identifies these columns. Triggers that do not specify a column list are not included in this view. Only those columns are shown that the current user owns or has some privilege other than SELECT on.

**Table 37.52. triggered_update_columns Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the trigger (always the current database)</td>
</tr>
<tr>
<td>trigger_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the trigger</td>
</tr>
<tr>
<td>trigger_name</td>
<td>sql_identifier</td>
<td>Name of the trigger</td>
</tr>
<tr>
<td>event_object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that the trigger is defined on (always the current database)</td>
</tr>
<tr>
<td>event_object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that the trigger is defined on</td>
</tr>
<tr>
<td>event_object_table</td>
<td>sql_identifier</td>
<td>Name of the table that the trigger is defined on</td>
</tr>
<tr>
<td>event_object_column</td>
<td>sql_identifier</td>
<td>Name of the column that the trigger is defined on</td>
</tr>
</tbody>
</table>

37.55. triggers

The view `triggers` contains all triggers defined in the current database on tables and views that the current user owns or has some privilege other than SELECT on.

**Table 37.53. triggers Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the trigger (always the current database)</td>
</tr>
<tr>
<td>trigger_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the trigger</td>
</tr>
<tr>
<td>trigger_name</td>
<td>sql_identifier</td>
<td>Name of the trigger</td>
</tr>
<tr>
<td>event_manipulation</td>
<td>character_data</td>
<td>Event that fires the trigger (INSERT, UPDATE, or DELETE)</td>
</tr>
<tr>
<td>event_object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that the trigger is defined on (always the current database)</td>
</tr>
<tr>
<td>event_object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that the trigger is defined on</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------</td>
<td>--------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>event_object_table</td>
<td>sql_identifier</td>
<td>Name of the table that the trigger is defined on</td>
</tr>
<tr>
<td>action_order</td>
<td>cardinal_number</td>
<td>Firing order among triggers on the same table having the same event_manipulation, action_timing, and action_orientation. In PostgreSQL, triggers are fired in name order, so this column reflects that.</td>
</tr>
<tr>
<td>action_condition</td>
<td>character_data</td>
<td>WHEN condition of the trigger, null if none (also null if the table is not owned by a currently enabled role)</td>
</tr>
<tr>
<td>action_statement</td>
<td>character_data</td>
<td>Statement that is executed by the trigger (currently always EXECUTE FUNCTION function(...)</td>
</tr>
<tr>
<td>action_orientation</td>
<td>character_data</td>
<td>Identifies whether the trigger fires once for each processed row or once for each statement (ROW or STATEMENT)</td>
</tr>
<tr>
<td>action_timing</td>
<td>character_data</td>
<td>Time at which the trigger fires (BEFORE, AFTER, or INSTEAD OF)</td>
</tr>
<tr>
<td>action_reference_old_table</td>
<td>sql_identifier</td>
<td>Name of the “old” transition table, or null if none</td>
</tr>
<tr>
<td>action_reference_new_table</td>
<td>sql_identifier</td>
<td>Name of the “new” transition table, or null if none</td>
</tr>
<tr>
<td>action_reference_old_row</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>action_reference_new_row</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>created</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

Triggers in PostgreSQL have two incompatibilities with the SQL standard that affect the representation in the information schema. First, trigger names are local to each table in PostgreSQL, rather than being independent schema objects. Therefore there can be duplicate trigger names defined in one schema, so long as they belong to different tables. (trigger_catalog and trigger_schema are really the values pertaining to the table that the trigger is defined on.) Second, triggers can be defined to fire on multiple events in PostgreSQL (e.g., ON INSERT OR UPDATE), whereas the SQL standard only allows one. If a trigger is defined to fire on multiple events, it is represented as multiple rows in the information schema, one for each type of event. As a consequence of these two issues, the primary key of the view triggers is really (trigger_catalog, trigger_schema, event_object_table, trigger_name, event_manipulation) instead of (trigger_catalog, trigger_schema, trigger_name), which is what the SQL standard specifies. Nonetheless, if you define your triggers in a manner that conforms with the SQL standard (trigger names unique in the schema and only one event type per trigger), this will not affect you.
Note

Prior to PostgreSQL 9.1, this view’s columns action_timing, action_reference_old_table, action_reference_new_table, action_reference_old_row, and action_reference_new_row were named condition_timing, condition_reference_old_table, condition_reference_new_table, condition_reference_old_row, and condition_reference_new_row respectively. That was how they were named in the SQL:1999 standard. The new naming conforms to SQL:2003 and later.

37.56. udt_privileges

The view udt_privileges identifies USAGE privileges granted on user-defined types to a currently enabled role or by a currently enabled role. There is one row for each combination of type, grantor, and grantee. This view shows only composite types (see under Section 37.58 for why); see Section 37.57 for domain privileges.

Table 37.54. udt_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always TYPE USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

37.57. usage_privileges

The view usage_privileges identifies USAGE privileges granted on various kinds of objects to a currently enabled role or by a currently enabled role. In PostgreSQL, this currently applies to collations, domains, foreign-data wrappers, foreign servers, and sequences. There is one row for each combination of object, grantor, and grantee.

Since collations do not have real privileges in PostgreSQL, this view shows implicit non-grantable USAGE privileges granted by the owner to PUBLIC for all collations. The other object types, however, show real privileges.

In PostgreSQL, sequences also support SELECT and UPDATE privileges in addition to the USAGE privilege. These are nonstandard and therefore not visible in the information schema.

Table 37.55. usage_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
</tbody>
</table>
The view `user_defined_types` currently contains all composite types defined in the current database. Only those types are shown that the current user has access to (by way of being the owner or having some privilege).

SQL knows about two kinds of user-defined types: structured types (also known as composite types in PostgreSQL) and distinct types (not implemented in PostgreSQL). To be future-proof, use the column `user_defined_type_category` to differentiate between these. Other user-defined types such as base types and enums, which are PostgreSQL extensions, are not shown here. For domains, see Section 37.22 instead.

### Table 37.56. `user_defined_types` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>user_defined_type_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the type (always the current database)</td>
</tr>
<tr>
<td>user_defined_type_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the type</td>
</tr>
<tr>
<td>user_defined_type_name</td>
<td>sql_identifier</td>
<td>Name of the type</td>
</tr>
<tr>
<td>user_defined_type_category</td>
<td>character_data</td>
<td>Currently always STRUCTURED</td>
</tr>
<tr>
<td>is_instantiable</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_final</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>ordering_form</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>ordering_category</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------</td>
<td>------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>ordering_routine_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>ordering_routine_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>ordering_routine_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>reference_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>source_dtd_identifier</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>ref_dtd_identifier</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

37.59. **user_mapping_options**

The view `user_mapping_options` contains all the options defined for user mappings in the current database. Only those user mappings are shown where the current user has access to the corresponding foreign server (by way of being the owner or having some privilege).
### Table 37.57. user_mapping_options Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the user being mapped, or PUBLIC if the mapping is public</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server used by this mapping is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server used by this mapping</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option. This column will show as null unless the current user is the user being mapped, or the mapping is for PUBLIC and the current user is the server owner, or the current user is a superuser. The intent is to protect password information stored as user mapping option.</td>
</tr>
</tbody>
</table>

### 37.60. user_mappings

The view user_mappings contains all user mappings defined in the current database. Only those user mappings are shown where the current user has access to the corresponding foreign server (by way of being the owner or having some privilege).

### Table 37.58. user_mappings Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the user being mapped, or PUBLIC if the mapping is public</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server used by this mapping is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server used by this mapping</td>
</tr>
</tbody>
</table>

### 37.61. view_column_usage

The view view_column_usage identifies all columns that are used in the query expression of a view (the SELECT statement that defines the view). A column is only included if the table that contains the column is owned by a currently enabled role.

#### Note

Columns of system tables are not included. This should be fixed sometime.
Table 37.59. view_column_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>view_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the view (always the current database)</td>
</tr>
<tr>
<td>view_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the view</td>
</tr>
<tr>
<td>view_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column that is used by the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column that is used by the view</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column that is used by the view</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column that is used by the view</td>
</tr>
</tbody>
</table>

37.62. view_routine_usage

The view view_routine_usage identifies all routines (functions and procedures) that are used in the query expression of a view (the SELECT statement that defines the view). A routine is only included if that routine is owned by a currently enabled role.

Table 37.60. view_routine_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the view</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 37.40 for more information.</td>
</tr>
</tbody>
</table>

37.63. view_table_usage

The view view_table_usage identifies all tables that are used in the query expression of a view (the SELECT statement that defines the view). A table is only included if that table is owned by a currently enabled role.
Note
System tables are not included. This should be fixed sometime.

Table 37.61. view_table_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>view_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the view (always the current database)</td>
</tr>
<tr>
<td>view_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the view</td>
</tr>
<tr>
<td>view_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that is used by the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that is used by the view</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that is used by the view</td>
</tr>
</tbody>
</table>

37.64. views

The view views contains all views defined in the current database. Only those views are shown that the current user has access to (by way of being the owner or having some privilege).

Table 37.62. views Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the view</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>view_definition</td>
<td>character_data</td>
<td>Query expression defining the view (null if the view is not owned by a currently enabled role)</td>
</tr>
<tr>
<td>check_option</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_updatable</td>
<td>yes_or_no</td>
<td>YES if the view is updatable (allows UPDATE and DELETE), NO if not</td>
</tr>
<tr>
<td>is_insertable_into</td>
<td>yes_or_no</td>
<td>YES if the view is insertable into (allows INSERT), NO if not</td>
</tr>
<tr>
<td>is_trigger_updatable</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF UPDATE trigger defined on it, NO if not</td>
</tr>
</tbody>
</table>

1021
<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>is_trigger_deletable</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF DELETE trigger defined on it, NO if not</td>
</tr>
<tr>
<td>is_trigger_insertable_into</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF INSERT trigger defined on it, NO if not</td>
</tr>
</tbody>
</table>
Part V. Server Programming

This part is about extending the server functionality with user-defined functions, data types, triggers, etc. These are advanced topics which should probably be approached only after all the other user documentation about PostgreSQL has been understood. Later chapters in this part describe the server-side programming languages available in the PostgreSQL distribution as well as general issues concerning server-side programming languages. It is essential to read at least the earlier sections of Chapter 38 (covering functions) before diving into the material about server-side programming languages.
# Table of Contents

38. Extending SQL .............................................................. 1028

38.1. How Extensibility Works ........................................ 1028

38.2. The PostgreSQL Type System ........................................ 1028
  38.2.1. Base Types .................................................... 1028
  38.2.2. Container Types ............................................. 1028
  38.2.3. Domains ....................................................... 1029
  38.2.4. Pseudo-Types ................................................ 1029
  38.2.5. Polymorphic Types .......................................... 1029

38.3. User-defined Functions ........................................... 1030

38.4. User-defined Procedures .......................................... 1030

38.5. Query Language (SQL) Functions .................................. 1030
  38.5.1. Arguments for SQL Functions ................................ 1031
  38.5.2. SQL Functions on Base Types ................................. 1032
  38.5.3. SQL Functions on Composite Types .......................... 1034
  38.5.4. SQL Functions with Output Parameters ..................... 1037
  38.5.5. SQL Functions with Variable Numbers of Arguments ..... 1038
  38.5.6. SQL Functions with Default Values for Arguments ....... 1039
  38.5.7. SQL Functions as Table Sources ............................ 1040
  38.5.8. SQL Functions Returning Sets ............................... 1040
  38.5.9. SQL Functions Returning TABLE ............................. 1044
  38.5.10. Polymorphic SQL Functions ................................. 1044
  38.5.11. SQL Functions with Collations .............................. 1046

38.6. Function Overloading ................................................ 1046

38.7. Function Volatility Categories .................................... 1047

38.8. Procedural Language Functions .................................... 1049

38.9. Internal Functions .................................................. 1049

38.10. C-Language Functions ............................................. 1049
  38.10.1. Dynamic Loading ............................................ 1049
  38.10.2. Base Types in C-Language Functions ....................... 1051
  38.10.3. Version 1 Calling Conventions ............................. 1053
  38.10.4. Writing Code ................................................ 1056
  38.10.5. Compiling and Linking Dynamically-loaded Functions .... 1057
  38.10.6. Composite-type Arguments .................................. 1059
  38.10.7. Returning Rows (Composite Types) .......................... 1060
  38.10.8. Returning Sets .............................................. 1062
  38.10.9. Polymorphic Arguments and Return Types .................. 1067
  38.10.10. Transform Functions ....................................... 1069
  38.10.11. Shared Memory and LWLocks ............................... 1069
  38.10.12. Using C++ for Extensibility ............................... 1070

38.11. User-defined Aggregates ......................................... 1070
  38.11.1. Moving-Aggregate Mode .................................... 1072
  38.11.2. Polymorphic and Variadic Aggregates ...................... 1073
  38.11.3. Ordered-Set Aggregates ................................... 1075
  38.11.4. Partial Aggregation ........................................ 1076
  38.11.5. Support Functions for Aggregates ......................... 1077

38.12. User-defined Types ................................................ 1077
  38.12.1. TOAST Considerations ....................................... 1080

38.13. User-defined Operators ........................................... 1081

38.14. Operator Optimization Information ............................... 1082
  38.14.1. COMMUTATOR ................................................. 1082
  38.14.2. NEGATOR .................................................... 1083
  38.14.3. RESTRICT ................................................... 1083
  38.14.4. JOIN ......................................................... 1084
  38.14.5. HASHES ....................................................... 1085
  38.14.6. MERGES ..................................................... 1086
Chapter 38. Extending SQL

In the sections that follow, we will discuss how you can extend the PostgreSQL SQL query language by adding:

• functions (starting in Section 38.3)
• aggregates (starting in Section 38.11)
• data types (starting in Section 38.12)
• operators (starting in Section 38.13)
• operator classes for indexes (starting in Section 38.15)
• packages of related objects (starting in Section 38.16)

38.1. How Extensibility Works

PostgreSQL is extensible because its operation is catalog-driven. If you are familiar with standard relational database systems, you know that they store information about databases, tables, columns, etc., in what are commonly known as system catalogs. (Some systems call this the data dictionary.) The catalogs appear to the user as tables like any other, but the DBMS stores its internal bookkeeping in them. One key difference between PostgreSQL and standard relational database systems is that PostgreSQL stores much more information in its catalogs: not only information about tables and columns, but also information about data types, functions, access methods, and so on. These tables can be modified by the user, and since PostgreSQL bases its operation on these tables, this means that PostgreSQL can be extended by users. By comparison, conventional database systems can only be extended by changing hardcoded procedures in the source code or by loading modules specially written by the DBMS vendor.

The PostgreSQL server can moreover incorporate user-written code into itself through dynamic loading. That is, the user can specify an object code file (e.g., a shared library) that implements a new type or function, and PostgreSQL will load it as required. Code written in SQL is even more trivial to add to the server. This ability to modify its operation “on the fly” makes PostgreSQL uniquely suited for rapid prototyping of new applications and storage structures.

38.2. The PostgreSQL Type System

PostgreSQL data types can be divided into base types, container types, domains, and pseudo-types.

38.2.1. Base Types

Base types are those, like integer, that are implemented below the level of the SQL language (typically in a low-level language such as C). They generally correspond to what are often known as abstract data types. PostgreSQL can only operate on such types through functions provided by the user and only understands the behavior of such types to the extent that the user describes them. The built-in base types are described in Chapter 8.

Enumerated (enum) types can be considered as a subcategory of base types. The main difference is that they can be created using just SQL commands, without any low-level programming. Refer to Section 8.7 for more information.

38.2.2. Container Types

PostgreSQL has three kinds of “container” types, which are types that contain multiple values of other types. These are arrays, composites, and ranges.

Arrays can hold multiple values that are all of the same type. An array type is automatically created for each base type, composite type, range type, and domain type. But there are no arrays of arrays. So
far as the type system is concerned, multi-dimensional arrays are the same as one-dimensional arrays. Refer to Section 8.15 for more information.

Composite types, or row types, are created whenever the user creates a table. It is also possible to use CREATE TYPE to define a “stand-alone” composite type with no associated table. A composite type is simply a list of types with associated field names. A value of a composite type is a row or record of field values. Refer to Section 8.16 for more information.

A range type can hold two values of the same type, which are the lower and upper bounds of the range. Range types are user-created, although a few built-in ones exist. Refer to Section 8.17 for more information.

38.2.3. Domains

A domain is based on a particular underlying type and for many purposes is interchangeable with its underlying type. However, a domain can have constraints that restrict its valid values to a subset of what the underlying type would allow. Domains are created using the SQL command CREATE DOMAIN. Refer to Section 8.18 for more information.

38.2.4. Pseudo-Types

There are a few “pseudo-types” for special purposes. Pseudo-types cannot appear as columns of tables or components of container types, but they can be used to declare the argument and result types of functions. This provides a mechanism within the type system to identify special classes of functions. Table 8.25 lists the existing pseudo-types.

38.2.5. Polymorphic Types

Five pseudo-types of special interest are anyelement, anyarray, anynonarray, anyenum, and anyrange, which are collectively called polymorphic types. Any function declared using these types is said to be a polymorphic function. A polymorphic function can operate on many different data types, with the specific data type(s) being determined by the data types actually passed to it in a particular call.

Polymorphic arguments and results are tied to each other and are resolved to a specific data type when a query calling a polymorphic function is parsed. Each position (either argument or return value) declared as anyelement is allowed to have any specific actual data type, but in any given call they must all be the same actual type. Each position declared as anyarray can have any array data type, but similarly they must all be the same type. And similarly, positions declared as anyrange must all be the same range type. Furthermore, if there are positions declared anyarray and others declared anyelement, the actual array type in the anyarray positions must be an array whose elements are the same type appearing in the anyelement positions. Similarly, if there are positions declared anyrange and others declared anyelement, the actual range type in the anyrange positions must be a range whose subtype is the same type appearing in the anyelement positions. anynonarray is treated exactly the same as anyelement, but adds the additional constraint that the actual type must not be an array type. anyenum is treated exactly the same as anyelement, but adds the additional constraint that the actual type must be an enum type.

Thus, when more than one argument position is declared with a polymorphic type, the net effect is that only certain combinations of actual argument types are allowed. For example, a function declared as equal(anyelement, anyelement) will take any two input values, so long as they are of the same data type.

When the return value of a function is declared as a polymorphic type, there must be at least one argument position that is also polymorphic, and the actual data type supplied as the argument determines the actual result type for that call. For example, if there were not already an array subscripting mechanism, one could define a function that implements subcripting as subscript(anyarray,
integer) returns anyelement. This declaration constrains the actual first argument to be an array type, and allows the parser to infer the correct result type from the actual first argument’s type. Another example is that a function declared as \( f(\text{anyarray}) \) returns \( \text{anyenum} \) will only accept arrays of enum types.

Note that \( \text{anynonarray} \) and \( \text{anyenum} \) do not represent separate type variables; they are the same type as \( \text{anymelement} \), just with an additional constraint. For example, declaring a function as \( f(\text{anymelement}, \text{anyenum}) \) is equivalent to declaring it as \( f(\text{anyenum}, \text{anyenum}) \); both actual arguments have to be the same enum type.

A variadic function (one taking a variable number of arguments, as in Section 38.5.5) can be polymorphic: this is accomplished by declaring its last parameter as \( \text{VARIADIC anyarray} \). For purposes of argument matching and determining the actual result type, such a function behaves the same as if you had written the appropriate number of \( \text{anynonarray} \) parameters.

### 38.3. User-defined Functions

PostgreSQL provides four kinds of functions:

- query language functions (functions written in SQL) (Section 38.5)
- procedural language functions (functions written in, for example, PL/pgSQL or PL/Tcl) (Section 38.8)
- internal functions (Section 38.9)
- C-language functions (Section 38.10)

Every kind of function can take base types, composite types, or combinations of these as arguments (parameters). In addition, every kind of function can return a base type or a composite type. Functions can also be defined to return sets of base or composite values.

Many kinds of functions can take or return certain pseudo-types (such as polymorphic types), but the available facilities vary. Consult the description of each kind of function for more details.

It’s easiest to define SQL functions, so we’ll start by discussing those. Most of the concepts presented for SQL functions will carry over to the other types of functions.

Throughout this chapter, it can be useful to look at the reference page of the CREATE FUNCTION command to understand the examples better. Some examples from this chapter can be found in \texttt{funcs.sql} and \texttt{funcs.c} in the src/tutorial directory in the PostgreSQL source distribution.

### 38.4. User-defined Procedures

A procedure is a database object similar to a function. The difference is that a procedure does not return a value, so there is no return type declaration. While a function is called as part of a query or DML command, a procedure is called explicitly using the CALL statement.

The explanations on how to define user-defined functions in the rest of this chapter apply to procedures as well, except that the CREATE PROCEDURE command is used instead, there is no return type, and some other features such as strictness don’t apply.

Collectively, functions and procedures are also known as \textit{routines}. There are commands such as ALTER ROUTINE and DROP ROUTINE that can operate on functions and procedures without having to know which kind it is. Note, however, that there is no CREATE ROUTINE command.

### 38.5. Query Language (SQL) Functions
SQL functions execute an arbitrary list of SQL statements, returning the result of the last query in the list. In the simple (non-set) case, the first row of the last query's result will be returned. (Bear in mind that “the first row” of a multirow result is not well-defined unless you use ORDER BY.) If the last query happens to return no rows at all, the null value will be returned.

Alternatively, an SQL function can be declared to return a set (that is, multiple rows) by specifying the function's return type as SETOF sometype, or equivalently by declaring it as RETURNS TABLE(columns). In this case all rows of the last query's result are returned. Further details appear below.

The body of an SQL function must be a list of SQL statements separated by semicolons. A semicolon after the last statement is optional. Unless the function is declared to return void, the last statement must be a SELECT, or an INSERT, UPDATE, or DELETE that has a RETURNING clause.

Any collection of commands in the SQL language can be packaged together and defined as a function. Besides SELECT queries, the commands can include data modification queries (INSERT, UPDATE, and DELETE), as well as other SQL commands. (You cannot use transaction control commands, e.g. COMMIT, SAVEPOINT, and some utility commands, e.g. VACUUM, in SQL functions.) However, the final command must be a SELECT or have a RETURNING clause that returns whatever is specified as the function's return type. Alternatively, if you want to define a SQL function that performs actions but has no useful value to return, you can define it as returning void. For example, this function removes rows with negative salaries from the emp table:

```
CREATE FUNCTION clean_emp() RETURNS void AS 'DELETE FROM emp
WHERE salary < 0;
' LANGUAGE SQL;
SELECT clean_emp();
```

clean_emp
---------

(1 row)

**Note**

The entire body of a SQL function is parsed before any of it is executed. While a SQL function can contain commands that alter the system catalogs (e.g., CREATE TABLE), the effects of such commands will not be visible during parse analysis of later commands in the function. Thus, for example, CREATE TABLE foo (...); INSERT INTO foo VALUES (...); will not work as desired if packaged up into a single SQL function, since foo won't exist yet when the INSERT command is parsed. It's recommended to use PL/pgSQL instead of a SQL function in this type of situation.

The syntax of the CREATE FUNCTION command requires the function body to be written as a string constant. It is usually most convenient to use dollar quoting (see Section 4.1.2.4) for the string constant. If you choose to use regular single-quoted string constant syntax, you must double single quote marks ('') and backslashes (\) (assuming escape string syntax) in the body of the function (see Section 4.1.2.1).

### 38.5.1. Arguments for SQL Functions

Arguments of a SQL function can be referenced in the function body using either names or numbers. Examples of both methods appear below.
To use a name, declare the function argument as having a name, and then just write that name in the
function body. If the argument name is the same as any column name in the current SQL command
within the function, the column name will take precedence. To override this, qualify the argument
name with the name of the function itself, that is function_name.argument_name. (If this
would conflict with a qualified column name, again the column name wins. You can avoid the ambi-
guity by choosing a different alias for the table within the SQL command.)

In the older numeric approach, arguments are referenced using the syntax $n: $1 refers to the first
input argument, $2 to the second, and so on. This will work whether or not the particular argument
was declared with a name.

If an argument is of a composite type, then the dot notation, e.g., argname.fieldname or
$1.fieldname, can be used to access attributes of the argument. Again, you might need to qualify
the argument’s name with the function name to make the form with an argument name unambiguous.

SQL function arguments can only be used as data values, not as identifiers. Thus for example this
is reasonable:

INSERT INTO mytable VALUES ($1);

but this will not work:

INSERT INTO $1 VALUES (42);

Note

The ability to use names to reference SQL function arguments was added in Post-
greSQL 9.2. Functions to be used in older servers must use the $n notation.

38.5.2. SQL Functions on Base Types

The simplest possible SQL function has no arguments and simply returns a base type, such as in-
teger:

CREATE FUNCTION one() RETURNS integer AS $$
    SELECT 1 AS result;
$$ LANGUAGE SQL;

-- Alternative syntax for string literal:
CREATE FUNCTION one() RETURNS integer AS ' 
    SELECT 1 AS result;
' LANGUAGE SQL;

SELECT one();

    one
   ------
     1

Notice that we defined a column alias within the function body for the result of the function (with the
name result), but this column alias is not visible outside the function. Hence, the result is labeled
one instead of result.

It is almost as easy to define SQL functions that take base types as arguments:
CREATE FUNCTION add_em(x integer, y integer) RETURNS integer AS $$
SELECT x + y;
$$ LANGUAGE SQL;

SELECT add_em(1, 2) AS answer;

answer
-------
 3

Alternatively, we could dispense with names for the arguments and use numbers:

CREATE FUNCTION add_em(integer, integer) RETURNS integer AS $$
SELECT $1 + $2;
$$ LANGUAGE SQL;

SELECT add_em(1, 2) AS answer;

answer
-------
 3

Here is a more useful function, which might be used to debit a bank account:

CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS numeric AS $$
UPDATE bank
  SET balance = balance - debit
WHERE accountno = tf1.accountno;
SELECT 1;
$$ LANGUAGE SQL;

A user could execute this function to debit account 17 by $100.00 as follows:

SELECT tf1(17, 100.0);  

In this example, we chose the name accountno for the first argument, but this is the same as the name of a column in the bank table. Within the UPDATE command, accountno refers to the column bank.accountno, so tf1.accountno must be used to refer to the argument. We could of course avoid this by using a different name for the argument.

In practice one would probably like a more useful result from the function than a constant 1, so a more likely definition is:

CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS numeric AS $$
UPDATE bank
  SET balance = balance - debit
WHERE accountno = tf1.accountno;
SELECT balance FROM bank WHERE accountno = tf1.accountno;
$$ LANGUAGE SQL;

which adjusts the balance and returns the new balance. The same thing could be done in one command using RETURNING:
CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS numeric AS $$
    UPDATE bank
    SET balance = balance - debit
    WHERE accountno = tf1.accountno
    RETURNING balance;
$$ LANGUAGE SQL;

A SQL function must return exactly its declared result type. This may require inserting an explicit cast. For example, suppose we wanted the previous add_em function to return type float8 instead. This won't work:

CREATE FUNCTION add_em(integer, integer) RETURNS float8 AS $$
    SELECT $1 + $2;
$$ LANGUAGE SQL;

even though in other contexts PostgreSQL would be willing to insert an implicit cast to convert integer to float8. We need to write it as

CREATE FUNCTION add_em(integer, integer) RETURNS float8 AS $$
    SELECT ($1 + $2)::float8;
$$ LANGUAGE SQL;

### 38.5.3. SQL Functions on Composite Types

When writing functions with arguments of composite types, we must not only specify which argument we want but also the desired attribute (field) of that argument. For example, suppose that emp is a table containing employee data, and therefore also the name of the composite type of each row of the table. Here is a function double_salary that computes what someone's salary would be if it were doubled:

CREATE TABLE emp (  
    name text,  
    salary numeric,  
    age integer,  
    cubicle point  
);

INSERT INTO emp VALUES ('Bill', 4200, 45, '(2,1)');

CREATE FUNCTION double_salary(emp) RETURNS numeric AS $$
    SELECT $1.salary * 2 AS salary;
$$ LANGUAGE SQL;

SELECT name, double_salary(emp.*) AS dream  
FROM emp  
WHERE emp.cubicle ~= point '(2,1)';

<table>
<thead>
<tr>
<th>name</th>
<th>dream</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bill</td>
<td>8400</td>
</tr>
</tbody>
</table>

Notice the use of the syntax $1.salary to select one field of the argument row value. Also notice how the calling SELECT command uses table_name.* to select the entire current row of a table as a composite value. The table row can alternatively be referenced using just the table name, like this:
SELECT name, double_salary(emp) AS dream
FROM emp
WHERE emp.cubicle ~= point '(2,1)';

but this usage is deprecated since it's easy to get confused. (See Section 8.16.5 for details about these
two notations for the composite value of a table row.)

Sometimes it is handy to construct a composite argument value on-the-fly. This can be done with the
ROW construct. For example, we could adjust the data being passed to the function:

SELECT name, double_salary(ROW(name, salary*1.1, age, cubicle)) AS
dream
FROM emp;

It is also possible to build a function that returns a composite type. This is an example of a function
that returns a single emp row:

CREATE FUNCTION new_emp() RETURNS emp AS $$
SELECT text 'None' AS name,
       1000.0 AS salary,
       25 AS age,
       point '(2,2)' AS cubicle;
$$ LANGUAGE SQL;

In this example we have specified each of the attributes with a constant value, but any computation
could have been substituted for these constants.

Note two important things about defining the function:

• The select list order in the query must be exactly the same as that in which the columns appear in
  the table associated with the composite type. (Naming the columns, as we did above, is irrelevant
to the system.)

• We must ensure each expression's type matches the corresponding column of the composite type,
  inserting a cast if necessary. Otherwise we'll get errors like this:

    ERROR:  function declared to return emp returns varchar instead
            of text at column 1

    As with the base-type case, the function will not insert any casts automatically.

A different way to define the same function is:

CREATE FUNCTION new_emp() RETURNS emp AS $$
SELECT ROW('None', 1000.0, 25, '(2,2)')::emp;
$$ LANGUAGE SQL;

Here we wrote a SELECT that returns just a single column of the correct composite type. This isn't
really better in this situation, but it is a handy alternative in some cases — for example, if we need
to compute the result by calling another function that returns the desired composite value. Another
example is that if we are trying to write a function that returns a domain over composite, rather than a
plain composite type, it is always necessary to write it as returning a single column, since there is no
other way to produce a value that is exactly of the domain type.
Extending SQL

We could call this function directly either by using it in a value expression:

```
SELECT new_emp();
```

```
new_emp
--------------------------
(None,1000.0,25,"(2,2)"
```

or by calling it as a table function:

```
SELECT * FROM new_emp();
```

```
name | salary | age | cubicle
------|--------+-----+---------
None | 1000.0 | 25 | (2,2)
```

The second way is described more fully in Section 38.5.7.

When you use a function that returns a composite type, you might want only one field (attribute) from its result. You can do that with syntax like this:

```
SELECT (new_emp()).name;
```

```
name
------
None
```

The extra parentheses are needed to keep the parser from getting confused. If you try to do it without them, you get something like this:

```
SELECT new_emp().name;
ERROR:  syntax error at or near "."
LINE 1: SELECT new_emp().name;
    ^
```

Another option is to use functional notation for extracting an attribute:

```
SELECT name(new_emp());
```

```
name
------
None
```

As explained in Section 8.16.5, the field notation and functional notation are equivalent.

Another way to use a function returning a composite type is to pass the result to another function that accepts the correct row type as input:

```
CREATE FUNCTION getname(emp) RETURNS text AS $$
  SELECT $1.name;
$$ LANGUAGE SQL;
```

```
SELECT getname(new_emp());
```
38.5.4. SQL Functions with Output Parameters

An alternative way of describing a function's results is to define it with output parameters, as in this example:

```sql
CREATE FUNCTION add_em (IN x int, IN y int, OUT sum int)
AS 'SELECT x + y'
LANGUAGE SQL;

SELECT add_em(3,7);
add_em
--------
  10
(1 row)
```

This is not essentially different from the version of `add_em` shown in Section 38.5.2. The real value of output parameters is that they provide a convenient way of defining functions that return several columns. For example,

```sql
CREATE FUNCTION sum_n_product (x int, y int, OUT sum int, OUT product int)
AS 'SELECT x + y, x * y'
LANGUAGE SQL;

SELECT * FROM sum_n_product(11,42);
sum | product
-----+---------
 53 | 462
(1 row)
```

What has essentially happened here is that we have created an anonymous composite type for the result of the function. The above example has the same end result as

```sql
CREATE TYPE sum_prod AS (sum int, product int);

CREATE FUNCTION sum_n_product (int, int) RETURNS sum_prod
AS 'SELECT $1 + $2, $1 * $2'
LANGUAGE SQL;
```

but not having to bother with the separate composite type definition is often handy. Notice that the names attached to the output parameters are not just decoration, but determine the column names of the anonymous composite type. (If you omit a name for an output parameter, the system will choose a name on its own.)

Notice that output parameters are not included in the calling argument list when invoking such a function from SQL. This is because PostgreSQL considers only the input parameters to define the function's calling signature. That means also that only the input parameters matter when referencing the function for purposes such as dropping it. We could drop the above function with either of
DROP FUNCTION sum_n_product (x int, y int, OUT sum int, OUT product int);
DROP FUNCTION sum_n_product (int, int);

Parameters can be marked as IN (the default), OUT, INOUT, or VARIADIC. An INOUT parameter serves as both an input parameter (part of the calling argument list) and an output parameter (part of the result record type). VARIADIC parameters are input parameters, but are treated specially as described next.

### 38.5.5. SQL Functions with Variable Numbers of Arguments

SQL functions can be declared to accept variable numbers of arguments, so long as all the “optional” arguments are of the same data type. The optional arguments will be passed to the function as an array. The function is declared by marking the last parameter as VARIADIC; this parameter must be declared as being of an array type. For example:

```
CREATE FUNCTION mleast(VARIADIC arr numeric[]) RETURNS numeric AS $$
    SELECT min($1[i]) FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;
```

```
SELECT mleast(10, -1, 5, 4.4);
mleast
--------
-1
(1 row)
```

Effectively, all the actual arguments at or beyond the VARIADIC position are gathered up into a one-dimensional array, as if you had written

```
SELECT mleast(ARRAY[10, -1, 5, 4.4]);    -- doesn't work
```

You can’t actually write that, though — or at least, it will not match this function definition. A parameter marked VARIADIC matches one or more occurrences of its element type, not of its own type.

Sometimes it is useful to be able to pass an already-constructed array to a variadic function; this is particularly handy when one variadic function wants to pass on its array parameter to another one. Also, this is the only secure way to call a variadic function found in a schema that permits untrusted users to create objects; see Section 10.3. You can do this by specifying VARIADIC in the call:

```
SELECT mleast(VARIADIC ARRAY[10, -1, 5, 4.4]);
```

This prevents expansion of the function’s variadic parameter into its element type, thereby allowing the array argument value to match normally. VARIADIC can only be attached to the last actual argument of a function call.

Specifying VARIADIC in the call is also the only way to pass an empty array to a variadic function, for example:

```
SELECT mleast(VARIADIC ARRAY[]::numeric[]);
```

Simply writing `SELECT mleast()` does not work because a variadic parameter must match at least one actual argument. (You could define a second function also named `mleast`, with no parameters, if you wanted to allow such calls.)
The array element parameters generated from a variadic parameter are treated as not having any names of their own. This means it is not possible to call a variadic function using named arguments (Section 4.3), except when you specify VARIADIC. For example, this will work:

```sql
SELECT mleast(VARIADIC arr => ARRAY[10, -1, 5, 4.4]);
```

but not these:

```sql
SELECT mleast(arr => 10);
SELECT mleast(arr => ARRAY[10, -1, 5, 4.4]);
```

### 38.5.6. SQL Functions with Default Values for Arguments

Functions can be declared with default values for some or all input arguments. The default values are inserted whenever the function is called with insufficiently many actual arguments. Since arguments can only be omitted from the end of the actual argument list, all parameters after a parameter with a default value have to have default values as well. (Although the use of named argument notation could allow this restriction to be relaxed, it’s still enforced so that positional argument notation works sensibly.) Whether or not you use it, this capability creates a need for precautions when calling functions in databases where some users mistrust other users; see Section 10.3.

For example:

```sql
CREATE FUNCTION foo(a int, b int DEFAULT 2, c int DEFAULT 3)
RETURNS int
LANGUAGE SQL
AS $$
    SELECT $1 + $2 + $3;
$$;

SELECT foo(10, 20, 30);
foo
-----
  60
(1 row)

SELECT foo(10, 20);
foo
-----
  33
(1 row)

SELECT foo(10);
foo
-----
  15
(1 row)

SELECT foo();  -- fails since there is no default for the first argument
ERROR:  function foo() does not exist

The = sign can also be used in place of the key word DEFAULT.
38.5.7. SQL Functions as Table Sources

All SQL functions can be used in the FROM clause of a query, but it is particularly useful for functions returning composite types. If the function is defined to return a base type, the table function produces a one-column table. If the function is defined to return a composite type, the table function produces a column for each attribute of the composite type.

Here is an example:

```sql
CREATE TABLE foo (fooid int, foosubid int, fooname text);
INSERT INTO foo VALUES (1, 1, 'Joe');
INSERT INTO foo VALUES (1, 2, 'Ed');
INSERT INTO foo VALUES (2, 1, 'Mary');

CREATE FUNCTION getfoo(int) RETURNS foo AS $$
    SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;

SELECT *, upper(fooname) FROM getfoo(1) AS t1;
```

<table>
<thead>
<tr>
<th>fooid</th>
<th>foosubid</th>
<th>fooname</th>
<th>upper</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Joe</td>
<td>JOE</td>
</tr>
</tbody>
</table>

As the example shows, we can work with the columns of the function's result just the same as if they were columns of a regular table.

Note that we only got one row out of the function. This is because we did not use SETOF. That is described in the next section.

38.5.8. SQL Functions Returning Sets

When an SQL function is declared as returning SETOF sometype, the function's final query is executed to completion, and each row it outputs is returned as an element of the result set.

This feature is normally used when calling the function in the FROM clause. In this case each row returned by the function becomes a row of the table seen by the query. For example, assume that table foo has the same contents as above, and we say:

```sql
CREATE FUNCTION getfoo(int) RETURNS SETOF foo AS $$
    SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;

SELECT * FROM getfoo(1) AS t1;
```

Then we would get:

<table>
<thead>
<tr>
<th>fooid</th>
<th>foosubid</th>
<th>fooname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Joe</td>
</tr>
<tr>
<td>1</td>
<td>2</td>
<td>Ed</td>
</tr>
</tbody>
</table>

It is also possible to return multiple rows with the columns defined by output parameters, like this:
CREATE TABLE tab (y int, z int);
INSERT INTO tab VALUES (1, 2), (3, 4), (5, 6), (7, 8);

CREATE FUNCTION sum_n_product_with_tab (x int, OUT sum int, OUT product int)
RETURNS SETOF record
AS $$
SELECT $1 + tab.y, $1 * tab.y FROM tab;
$$ LANGUAGE SQL;

SELECT * FROM sum_n_product_with_tab(10);

<table>
<thead>
<tr>
<th>sum</th>
<th>product</th>
</tr>
</thead>
<tbody>
<tr>
<td>11</td>
<td>10</td>
</tr>
<tr>
<td>13</td>
<td>30</td>
</tr>
<tr>
<td>15</td>
<td>50</td>
</tr>
<tr>
<td>17</td>
<td>70</td>
</tr>
</tbody>
</table>

The key point here is that you must write RETURNS SETOF record to indicate that the function returns multiple rows instead of just one. If there is only one output parameter, write that parameter's type instead of record.

It is frequently useful to construct a query's result by invoking a set-returning function multiple times, with the parameters for each invocation coming from successive rows of a table or subquery. The preferred way to do this is to use the LATERAL key word, which is described in Section 7.2.1.5. Here is an example using a set-returning function to enumerate elements of a tree structure:

SELECT * FROM nodes;

<table>
<thead>
<tr>
<th>name</th>
<th>parent</th>
</tr>
</thead>
<tbody>
<tr>
<td>Top</td>
<td></td>
</tr>
<tr>
<td>Child1</td>
<td>Top</td>
</tr>
<tr>
<td>Child2</td>
<td>Top</td>
</tr>
<tr>
<td>Child3</td>
<td>Top</td>
</tr>
<tr>
<td>SubChild1</td>
<td>Child1</td>
</tr>
<tr>
<td>SubChild2</td>
<td>Child1</td>
</tr>
</tbody>
</table>

CREATE FUNCTION listchildren(text) RETURNS SETOF text AS $$
SELECT name FROM nodes WHERE parent = $1
$$ LANGUAGE SQL STABLE;

SELECT * FROM listchildren('Top');

<table>
<thead>
<tr>
<th>listchildren</th>
</tr>
</thead>
<tbody>
<tr>
<td>Child1</td>
</tr>
<tr>
<td>Child2</td>
</tr>
<tr>
<td>Child3</td>
</tr>
</tbody>
</table>

SELECT name, child FROM nodes, LATERAL listchildren(name) AS child;

<table>
<thead>
<tr>
<th>name</th>
<th>child</th>
</tr>
</thead>
<tbody>
<tr>
<td>Top</td>
<td>Child1</td>
</tr>
<tr>
<td>Top</td>
<td>Child2</td>
</tr>
<tr>
<td>Top</td>
<td>Child3</td>
</tr>
</tbody>
</table>
This example does not do anything that we couldn't have done with a simple join, but in more complex calculations the option to put some of the work into a function can be quite convenient.

Functions returning sets can also be called in the select list of a query. For each row that the query generates by itself, the set-returning function is invoked, and an output row is generated for each element of the function's result set. The previous example could also be done with queries like these:

```sql
SELECT listchildren('Top');
```

```output
listchildren
------------
Child1
Child2
Child3
(3 rows)
```

```sql
SELECT name, listchildren(name) FROM nodes;
```

```output
name | listchildren
--------+--------------
Top    | Child1
Top    | Child2
Top    | Child3
Child1 | SubChild1
Child1 | SubChild2
(5 rows)
```

In the last `SELECT`, notice that no output row appears for Child2, Child3, etc. This happens because `listchildren` returns an empty set for those arguments, so no result rows are generated. This is the same behavior as we got from an inner join to the function result when using the `LATERAL` syntax.

`PostgreSQL`'s behavior for a set-returning function in a query's select list is almost exactly the same as if the set-returning function had been written in a `LATERAL` FROM-clause item instead. For example,

```sql
SELECT x, generate_series(1,5) AS g FROM tab;
```

is almost equivalent to

```sql
SELECT x, g FROM tab, LATERAL generate_series(1,5) AS g;
```

It would be exactly the same, except that in this specific example, the planner could choose to put `g` on the outside of the nestloop join, since `g` has no actual lateral dependency on `tab`. That would result in a different output row order. Set-returning functions in the select list are always evaluated as though they are on the inside of a nestloop join with the rest of the FROM clause, so that the function(s) are run to completion before the next row from the FROM clause is considered.

If there is more than one set-returning function in the query's select list, the behavior is similar to what you get from putting the functions into a single `LATERAL ROWS FROM( ... ) FROM-clause item. For each row from the underlying query, there is an output row using the first result from each function, then an output row using the second result, and so on. If some of the set-returning functions produce fewer outputs than others, null values are substituted for the missing data, so that the total number of rows emitted for one underlying row is the same as for the set-returning function that produced the most outputs. Thus the set-returning functions run “in lockstep” until they are all exhausted, and then execution continues with the next underlying row.
Set-returning functions can be nested in a select list, although that is not allowed in FROM-clause items. In such cases, each level of nesting is treated separately, as though it were a separate LATERAL ROWS FROM( ... ) item. For example, in

```
SELECT srf1(srf2(x), srf3(y)), srf4(srf5(z)) FROM tab;
```

the set-returning functions srf2, srf3, and srf5 would be run in lockstep for each row of tab, and then srf1 and srf4 would be applied in lockstep to each row produced by the lower functions.

Set-returning functions cannot be used within conditional-evaluation constructs, such as CASE or COALESCE. For example, consider

```
SELECT x, CASE WHEN x > 0 THEN generate_series(1, 5) ELSE 0 END FROM tab;
```

It might seem that this should produce five repetitions of input rows that have x > 0, and a single repetition of those that do not; but actually, because `generate_series(1, 5)` would be run in an implicit LATERAL FROM item before the CASE expression is ever evaluated, it would produce five repetitions of every input row. To reduce confusion, such cases produce a parse-time error instead.

**Note**

If a function's last command is INSERT, UPDATE, or DELETE with RETURNING, that command will always be executed to completion, even if the function is not declared with SETOF or the calling query does not fetch all the result rows. Any extra rows produced by the RETURNING clause are silently dropped, but the commanded table modifications still happen (and are all completed before returning from the function).

**Note**

Before PostgreSQL 10, putting more than one set-returning function in the same select list did not behave very sensibly unless they always produced equal numbers of rows. Otherwise, what you got was a number of output rows equal to the least common multiple of the numbers of rows produced by the set-returning functions. Also, nested set-returning functions did not work as described above; instead, a set-returning function could have at most one set-returning argument, and each nest of set-returning functions was run independently. Also, conditional execution (set-returning functions inside CASE etc) was previously allowed, complicating things even more. Use of the LATERAL syntax is recommended when writing queries that need to work in older PostgreSQL versions, because that will give consistent results across different versions. If you have a query that is relying on conditional execution of a set-returning function, you may be able to fix it by moving the conditional test into a custom set-returning function. For example,

```
SELECT x, CASE WHEN y > 0 THEN generate_series(1, z) ELSE 5 END FROM tab;
```

could become

```
CREATE FUNCTION case_generate_series(cond bool, start int, fin int, els int)
RETURNS SETOF int AS $$
```
BEGIN
    IF cond THEN
        RETURN QUERY SELECT generate_series(start, fin);
    ELSE
        RETURN QUERY SELECT els;
    END IF;
END$$ LANGUAGE plpgsql;

SELECT x, case_generate_series(y > 0, 1, z, 5) FROM tab;

This formulation will work the same in all versions of PostgreSQL.

38.5.9. SQL Functions Returning TABLE

There is another way to declare a function as returning a set, which is to use the syntax RETURNS TABLE(columns). This is equivalent to using one or more OUT parameters plus marking the function as returning SETOF record (or SETOF a single output parameter’s type, as appropriate). This notation is specified in recent versions of the SQL standard, and thus may be more portable than using SETOF.

For example, the preceding sum-and-product example could also be done this way:

CREATE FUNCTION sum_n_product_with_tab (x int)
RETURNS TABLE(sum int, product int) AS $$
SELECT $1 + tab.y, $1 * tab.y FROM tab;
$$ LANGUAGE SQL;

It is not allowed to use explicit OUT or INOUT parameters with the RETURNS TABLE notation — you must put all the output columns in the TABLE list.

38.5.10. Polymorphic SQL Functions

SQL functions can be declared to accept and return the polymorphic types anyelement, anyarray, anynonarray, anyenum, and anyrange. See Section 38.2.5 for a more detailed explanation of polymorphic functions. Here is a polymorphic function make_array that builds up an array from two arbitrary data type elements:

CREATE FUNCTION make_array(anyelement, anyelement) RETURNS anyarray
AS $$
SELECT ARRAY[$1, $2];
$$ LANGUAGE SQL;

SELECT make_array(1, 2) AS intarray, make_array('a'::text, 'b') AS textarray;

<table>
<thead>
<tr>
<th>intarray</th>
<th>textarray</th>
</tr>
</thead>
<tbody>
<tr>
<td>{1,2}</td>
<td>{a,b}</td>
</tr>
<tr>
<td>(1 row)</td>
<td></td>
</tr>
</tbody>
</table>

Notice the use of the typecast 'a'::text to specify that the argument is of type text. This is required if the argument is just a string literal, since otherwise it would be treated as type unknown, and array of unknown is not a valid type. Without the typecast, you will get errors like this:
ERROR: could not determine polymorphic type because input has type "unknown"

It is permitted to have polymorphic arguments with a fixed return type, but the converse is not. For example:

```sql
CREATE FUNCTION is_greater(anyelement, anyelement) RETURNS boolean
AS $$
    SELECT $1 > $2;
$$ LANGUAGE SQL;

SELECT is_greater(1, 2);
---------
f
(1 row)
```

```sql
CREATE FUNCTION invalid_func() RETURNS anyelement AS $$
    SELECT 1;
$$ LANGUAGE SQL;
ERROR: cannot determine result data type
DETAIL: A function returning a polymorphic type must have at least one polymorphic argument.
```

Polymorphism can be used with functions that have output arguments. For example:

```sql
CREATE FUNCTION dup (f1 anyelement, OUT f2 anyelement, OUT f3 anyarray)
AS 'select $1, array[$1,$1]' LANGUAGE SQL;

SELECT * FROM dup(22);
f2 | f3
---+---------
22 | {22,22}
(1 row)
```

Polymorphism can also be used with variadic functions. For example:

```sql
CREATE FUNCTION anyleast (VARIADIC anyarray) RETURNS anyelement AS $$
    SELECT min($1[i]) FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;

SELECT anyleast(10, -1, 5, 4);
---------
-1
(1 row)

SELECT anyleast('abc'::text, 'def');
---------
abc
(1 row)
```
CREATE FUNCTION concat_values(text, VARIADIC anyarray) RETURNS text
AS $$
SELECT array_to_string($2, $1);
$$ LANGUAGE SQL;

SELECT concat_values('|', 1, 4, 2);
concat_values
---------------
1|4|2
(1 row)

38.5.11. SQL Functions with Collations

When a SQL function has one or more parameters of collatable data types, a collation is identified for each function call depending on the collations assigned to the actual arguments, as described in Section 23.2. If a collation is successfully identified (i.e., there are no conflicts of implicit collations among the arguments) then all the collatable parameters are treated as having that collation implicitly. This will affect the behavior of collation-sensitive operations within the function. For example, using the anyleast function described above, the result of

SELECT anyleast('abc'::text, 'ABC');

will depend on the database's default collation. In C locale the result will be ABC, but in many other locales it will be abc. The collation to use can be forced by adding a COLLATE clause to any of the arguments, for example

SELECT anyleast('abc'::text, 'ABC' COLLATE "C");

Alternatively, if you wish a function to operate with a particular collation regardless of what it is called with, insert COLLATE clauses as needed in the function definition. This version of anyleast would always use en_US locale to compare strings:

CREATE FUNCTION anyleast (VARIADIC anyarray) RETURNS anyelement AS $$
SELECT min($1[i] COLLATE "en_US") FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;

But note that this will throw an error if applied to a non-collatable data type.

If no common collation can be identified among the actual arguments, then a SQL function treats its parameters as having their data types' default collation (which is usually the database's default collation, but could be different for parameters of domain types).

The behavior of collatable parameters can be thought of as a limited form of polymorphism, applicable only to textual data types.

38.6. Function Overloading

More than one function can be defined with the same SQL name, so long as the arguments they take are different. In other words, function names can be overloaded. Whether or not you use it, this capability entails security precautions when calling functions in databases where some users mistrust other users; see Section 10.3. When a query is executed, the server will determine which function to call from the data types and the number of the provided arguments. Overloading can also be used to simulate functions with a variable number of arguments, up to a finite maximum number.
When creating a family of overloaded functions, one should be careful not to create ambiguities. For instance, given the functions:

```sql
CREATE FUNCTION test(int, real) RETURNS ...
CREATE FUNCTION test(smallint, double precision) RETURNS ...
```

it is not immediately clear which function would be called with some trivial input like `test(1, 1.5)`. The currently implemented resolution rules are described in Chapter 10, but it is unwise to design a system that subtly relies on this behavior.

A function that takes a single argument of a composite type should generally not have the same name as any attribute (field) of that type. Recall that `attribute(table)` is considered equivalent to `table.attribute`. In the case that there is an ambiguity between a function on a composite type and an attribute of the composite type, the attribute will always be used. It is possible to override that choice by schema-qualifying the function name (that is, `schema.func(table)`), but it's better to avoid the problem by not choosing conflicting names.

Another possible conflict is between variadic and non-variadic functions. For instance, it is possible to create both `foo(numeric)` and `foo(VARIADIC numeric[])`. In this case it is unclear which one should be matched to a call providing a single numeric argument, such as `foo(10.1)`. The rule is that the function appearing earlier in the search path is used, or if the two functions are in the same schema, the non-variadic one is preferred.

When overloading C-language functions, there is an additional constraint: The C name of each function in the family of overloaded functions must be different from the C names of all other functions, either internal or dynamically loaded. If this rule is violated, the behavior is not portable. You might get a run-time linker error, or one of the functions will get called (usually the internal one). The alternative form of the AS clause for the SQL `CREATE FUNCTION` command decouples the SQL function name from the function name in the C source code. For instance:

```sql
CREATE FUNCTION test(int) RETURNS int
    AS 'filename', 'test_1arg'
    LANGUAGE C;
CREATE FUNCTION test(int, int) RETURNS int
    AS 'filename', 'test_2arg'
    LANGUAGE C;
```

The names of the C functions here reflect one of many possible conventions.

### 38.7. Function Volatility Categories

Every function has a volatility classification, with the possibilities being VOLATILE, STABLE, or IMMUTABLE. VOLATILE is the default if the `CREATE FUNCTION` command does not specify a category. The volatility category is a promise to the optimizer about the behavior of the function:

- A VOLATILE function can do anything, including modifying the database. It can return different results on successive calls with the same arguments. The optimizer makes no assumptions about the behavior of such functions. A query using a volatile function will re-evaluate the function at every row where its value is needed.

- A STABLE function cannot modify the database and is guaranteed to return the same results given the same arguments for all rows within a single statement. This category allows the optimizer to optimize multiple calls of the function to a single call. In particular, it is safe to use an expression containing such a function in an index scan condition. (Since an index scan will evaluate the comparison value only once, not once at each row, it is not valid to use a VOLATILE function in an index scan condition.)
An **IMMUTABLE** function cannot modify the database and is guaranteed to return the same results given the same arguments forever. This category allows the optimizer to pre-evaluate the function when a query calls it with constant arguments. For example, a query like `SELECT ... WHERE x = 2 + 2` can be simplified on sight to `SELECT ... WHERE x = 4`, because the function underlying the integer addition operator is marked **IMMUTABLE**.

For best optimization results, you should label your functions with the strictest volatility category that is valid for them.

Any function with side-effects **must** be labeled **VOLATILE**, so that calls to it cannot be optimized away. Even a function with no side-effects needs to be labeled **VOLATILE** if its value can change within a single query; some examples are `random()`, `currval()`, `timeofday()`.

Another important example is that the `current_timestamp` family of functions qualify as **STABLE**, since their values do not change within a transaction.

There is relatively little difference between **STABLE** and **IMMUTABLE** categories when considering simple interactive queries that are planned and immediately executed: it doesn't matter a lot whether a function is executed once during planning or once during query execution startup. But there is a big difference if the plan is saved and reused later. Labeling a function **IMMUTABLE** when it really isn't might allow it to be prematurely folded to a constant during planning, resulting in a stale value being re-used during subsequent uses of the plan. This is a hazard when using prepared statements or when using function languages that cache plans (such as PL/pgSQL).

For functions written in SQL or in any of the standard procedural languages, there is a second important property determined by the volatility category, namely the visibility of any data changes that have been made by the SQL command that is calling the function. A **VOLATILE** function will see such changes, a **STABLE** or **IMMUTABLE** function will not. This behavior is implemented using the snapshotting behavior of MVCC (see Chapter 13): **STABLE** and **IMMUTABLE** functions use a snapshot established as of the start of the calling query, whereas **VOLATILE** functions obtain a fresh snapshot at the start of each query they execute.

**Note**

Functions written in C can manage snapshots however they want, but it's usually a good idea to make C functions work this way too.

Because of this snapshotting behavior, a function containing only `SELECT` commands can safely be marked **STABLE**, even if it selects from tables that might be undergoing modifications by concurrent queries. PostgreSQL will execute all commands of a **STABLE** function using the snapshot established for the calling query, and so it will see a fixed view of the database throughout that query.

The same snapshotting behavior is used for `SELECT` commands within **IMMUTABLE** functions. It is generally unwise to select from database tables within an **IMMUTABLE** function at all, since the immutability will be broken if the table contents ever change. However, PostgreSQL does not enforce that you do not do that.

A common error is to label a function **IMMUTABLE** when its results depend on a configuration parameter. For example, a function that manipulates timestamps might well have results that depend on the TimeZone setting. For safety, such functions should be labeled **STABLE** instead.

**Note**

PostgreSQL requires that **STABLE** and **IMMUTABLE** functions contain no SQL commands other than `SELECT` to prevent data modification. (This is not a completely bulletproof test, since such functions could still call **VOLATILE** functions that modify the database. If you do that, you will find that the **STABLE** or **IMMUTABLE** function does...
38.8. Procedural Language Functions

PostgreSQL allows user-defined functions to be written in other languages besides SQL and C. These other languages are generically called *procedural languages* (PLs). Procedural languages aren't built into the PostgreSQL server; they are offered by loadable modules. See Chapter 42 and following chapters for more information.

38.9. Internal Functions

Internal functions are functions written in C that have been statically linked into the PostgreSQL server. The “body” of the function definition specifies the C-language name of the function, which need not be the same as the name being declared for SQL use. (For reasons of backward compatibility, an empty body is accepted as meaning that the C-language function name is the same as the SQL name.)

Normally, all internal functions present in the server are declared during the initialization of the database cluster (see Section 18.2), but a user could use `CREATE FUNCTION` to create additional alias names for an internal function. Internal functions are declared in `CREATE FUNCTION` with language name `internal`. For instance, to create an alias for the `sqrt` function:

```sql
CREATE FUNCTION square_root(double precision) RETURNS double precision AS 'dsqrt'
    LANGUAGE internal
    STRICT;
```

(Most internal functions expect to be declared “strict”.)

Note

Not all “predefined” functions are “internal” in the above sense. Some predefined functions are written in SQL.

38.10. C-Language Functions

User-defined functions can be written in C (or a language that can be made compatible with C, such as C++). Such functions are compiled into dynamically loadable objects (also called shared libraries) and are loaded by the server on demand. The dynamic loading feature is what distinguishes “C language” functions from “internal” functions — the actual coding conventions are essentially the same for both. (Hence, the standard internal function library is a rich source of coding examples for user-defined C functions.)

Currently only one calling convention is used for C functions (“version 1”). Support for that calling convention is indicated by writing a `PG_FUNCTION_INFO_V1()` macro call for the function, as illustrated below.

38.10.1. Dynamic Loading

The first time a user-defined function in a particular loadable object file is called in a session, the dynamic loader loads that object file into memory so that the function can be called. The `CREATE`
Extending SQL

The following algorithm is used to locate the shared object file based on the name given in the `CREATE FUNCTION` command:

1. If the name is an absolute path, the given file is loaded.
2. If the name starts with the string `$libdir`, that part is replaced by the PostgreSQL package library directory name, which is determined at build time.
3. If the name does not contain a directory part, the file is searched for in the path specified by the configuration variable `dynamic_library_path`.
4. Otherwise (the file was not found in the path, or it contains a non-absolute directory part), the dynamic loader will try to take the name as given, which will most likely fail. (It is unreliable to depend on the current working directory.)

If this sequence does not work, the platform-specific shared library file name extension (often `.so`) is appended to the given name and this sequence is tried again. If that fails as well, the load will fail.

It is recommended to locate shared libraries either relative to `$libdir` or through the dynamic library path. This simplifies version upgrades if the new installation is at a different location. The actual directory that `$libdir` stands for can be found out with the command `pg_config --pkglibdir`.

The user ID the PostgreSQL server runs as must be able to traverse the path to the file you intend to load. Making the file or a higher-level directory not readable and/or not executable by the postgres user is a common mistake.

In any case, the file name that is given in the `CREATE FUNCTION` command is recorded literally in the system catalogs, so if the file needs to be loaded again the same procedure is applied.

Note

PostgreSQL will not compile a C function automatically. The object file must be compiled before it is referenced in a `CREATE FUNCTION` command. See Section 38.10.5 for additional information.

To ensure that a dynamically loaded object file is not loaded into an incompatible server, PostgreSQL checks that the file contains a “magic block” with the appropriate contents. This allows the server to detect obvious incompatibilities, such as code compiled for a different major version of PostgreSQL.

To include a magic block, write this in one (and only one) of the module source files, after having included the header `fmgr.h`:

```
PG_MODULE_MAGIC;
```

After it is used for the first time, a dynamically loaded object file is retained in memory. Future calls in the same session to the function(s) in that file will only incur the small overhead of a symbol table lookup. If you need to force a reload of an object file, for example after recompiling it, begin a fresh session.

Optionally, a dynamically loaded file can contain initialization and finalization functions. If the file includes a function named `_PG_init`, that function will be called immediately after loading the file. The function receives no parameters and should return void. If the file includes a function named `_PG_fini`, that function will be called immediately before unloading the file. Likewise, the function...
receives no parameters and should return void. Note that _PG_fini will only be called during an unload of the file, not during process termination. (Presently, unloads are disabled and will never occur, but this may change in the future.)

38.10.2. Base Types in C-Language Functions

To know how to write C-language functions, you need to know how PostgreSQL internally represents base data types and how they can be passed to and from functions. Internally, PostgreSQL regards a base type as a “blob of memory”. The user-defined functions that you define over a type in turn define the way that PostgreSQL can operate on it. That is, PostgreSQL will only store and retrieve the data from disk and use your user-defined functions to input, process, and output the data.

Base types can have one of three internal formats:

- pass by value, fixed-length
- pass by reference, fixed-length
- pass by reference, variable-length

By-value types can only be 1, 2, or 4 bytes in length (also 8 bytes, if sizeof(Datum) is 8 on your machine). You should be careful to define your types such that they will be the same size (in bytes) on all architectures. For example, the long type is dangerous because it is 4 bytes on some machines and 8 bytes on others, whereas int type is 4 bytes on most Unix machines. A reasonable implementation of the int4 type on Unix machines might be:

```c
/* 4-byte integer, passed by value */
typedef int int4;
```

(The actual PostgreSQL C code calls this type int32, because it is a convention in C that intXX means XX bits. Note therefore also that the C type int8 is 1 byte in size. The SQL type int8 is called int64 in C. See also Table 38.1.)

On the other hand, fixed-length types of any size can be passed by-reference. For example, here is a sample implementation of a PostgreSQL type:

```c
/* 16-byte structure, passed by reference */
typedef struct
{
    double x, y;
} Point;
```

Only pointers to such types can be used when passing them in and out of PostgreSQL functions. To return a value of such a type, allocate the right amount of memory with palloc, fill in the allocated memory, and return a pointer to it. (Also, if you just want to return the same value as one of your input arguments that's of the same data type, you can skip the extra palloc and just return the pointer to the input value.)

Finally, all variable-length types must also be passed by reference. All variable-length types must begin with an opaque length field of exactly 4 bytes, which will be set by SET_VARSIZE; never set this field directly! All data to be stored within that type must be located in the memory immediately following that length field. The length field contains the total length of the structure, that is, it includes the size of the length field itself.

Another important point is to avoid leaving any uninitialized bits within data type values; for example, take care to zero out any alignment padding bytes that might be present in structs. Without this, logically-equivalent constants of your data type might be seen as unequal by the planner, leading to inefficient (though not incorrect) plans.
As an example, we can define the type `text` as follows:

```c
typedef struct {
    int32 length;
    char data[FLEXIBLE_ARRAY_MEMBER];
} text;
```

The `[FLEXIBLE_ARRAY_MEMBER]` notation means that the actual length of the data part is not specified by this declaration.

When manipulating variable-length types, we must be careful to allocate the correct amount of memory and set the length field correctly. For example, if we wanted to store 40 bytes in a `text` structure, we might use a code fragment like this:

```c
#include "postgres.h"
...
char buffer[40]; /* our source data */
...
text *destination = (text *) palloc(VARHDRSZ + 40);
SET_VARSIZE(destination, VARHDRSZ + 40);
memcpy(destination->data, buffer, 40);
...
```

`VARHDRSZ` is the same as `sizeof(int32)`, but it's considered good style to use the macro `VARHDRSZ` to refer to the size of the overhead for a variable-length type. Also, the length field must be set using the `SET_VARSIZE` macro, not by simple assignment.

Table 38.1 specifies which C type corresponds to which SQL type when writing a C-language function that uses a built-in type of PostgreSQL. The “Defined In” column gives the header file that needs to be included to get the type definition. (The actual definition might be in a different file that is included by the listed file. It is recommended that users stick to the defined interface.) Note that you should always include `postgres.h` first in any source file, because it declares a number of things that you will need anyway.

### Table 38.1. Equivalent C Types for Built-in SQL Types

<table>
<thead>
<tr>
<th>SQL Type</th>
<th>C Type</th>
<th>Defined In</th>
</tr>
</thead>
<tbody>
<tr>
<td>abstime</td>
<td>AbsoluteTime</td>
<td><code>utils/nabstime.h</code></td>
</tr>
<tr>
<td>bigint (int8)</td>
<td>int64</td>
<td><code>postgres.h</code></td>
</tr>
<tr>
<td>boolean</td>
<td>bool</td>
<td><code>postgres.h</code> (maybe compiler built-in)</td>
</tr>
<tr>
<td>box</td>
<td>BOX*</td>
<td><code>utils/geo_decls.h</code></td>
</tr>
<tr>
<td>bytea</td>
<td>bytea*</td>
<td><code>postgres.h</code></td>
</tr>
<tr>
<td>&quot;char&quot;</td>
<td>char</td>
<td>(compiler built-in)</td>
</tr>
<tr>
<td>character</td>
<td>BpChar*</td>
<td><code>postgres.h</code></td>
</tr>
</tbody>
</table>
### SQL Type

<table>
<thead>
<tr>
<th>SQL Type</th>
<th>C Type</th>
<th>Defined In</th>
</tr>
</thead>
<tbody>
<tr>
<td>cid</td>
<td>CommandId</td>
<td>postgres.h</td>
</tr>
<tr>
<td>date</td>
<td>DateADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>smallint(int2)</td>
<td>int16</td>
<td>postgres.h</td>
</tr>
<tr>
<td>int2vector</td>
<td>int2vector*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>integer(int4)</td>
<td>int32</td>
<td>postgres.h</td>
</tr>
<tr>
<td>real(float4)</td>
<td>float4*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>double precision (float8)</td>
<td>float8*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>interval</td>
<td>Interval*</td>
<td>datatype/timestamp.h</td>
</tr>
<tr>
<td>lseg</td>
<td>LSEG*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>name</td>
<td>Name</td>
<td>postgres.h</td>
</tr>
<tr>
<td>oid</td>
<td>Oid</td>
<td>postgres.h</td>
</tr>
<tr>
<td>oidvector</td>
<td>oidvector*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>path</td>
<td>PATH*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>point</td>
<td>POINT*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>regproc</td>
<td>regproc</td>
<td>postgres.h</td>
</tr>
<tr>
<td>reltime</td>
<td>RelativeTime</td>
<td>utils/nabstime.h</td>
</tr>
<tr>
<td>text</td>
<td>text*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>tid</td>
<td>TransactionId</td>
<td>storage/itemptr.h</td>
</tr>
<tr>
<td>time</td>
<td>TimeADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>time with time zone</td>
<td>TimeTzADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>timestamp</td>
<td>Timestamp*</td>
<td>datatype/timestamp.h</td>
</tr>
<tr>
<td>tinterval</td>
<td>TimeInterval</td>
<td>utils/nabstime.h</td>
</tr>
<tr>
<td>varchar</td>
<td>VarChar*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>xid</td>
<td>TransactionId</td>
<td>postgres.h</td>
</tr>
</tbody>
</table>

Now that we’ve gone over all of the possible structures for base types, we can show some examples of real functions.

### 38.10.3. Version 1 Calling Conventions

The version-1 calling convention relies on macros to suppress most of the complexity of passing arguments and results. The C declaration of a version-1 function is always:

```c
Datum funcname(PG_FUNCTION_ARGS)
```

In addition, the macro call:

```c
PG_FUNCTION_INFO_V1(funcname);
```

must appear in the same source file. (Conventionally, it’s written just before the function itself.) This macro call is not needed for internal-language functions, since PostgreSQL assumes that all internal functions use the version-1 convention. It is, however, required for dynamically-loaded functions.

In a version-1 function, each actual argument is fetched using a `PG_GETARG_XXX()` macro that corresponds to the argument’s data type. In non-strict functions there needs to be a previous check about ar-
argument null-ness using PG_ARGNULL_xxx(). The result is returned using a PG_RETURN_xxx() macro for the return type. PG_GETARG_xxx() takes as its argument the number of the function argument to fetch, where the count starts at 0. PG_RETURN_xxx() takes as its argument the actual value to return.

Here are some examples using the version-1 calling convention:

```
#include "postgres.h"
#include <string.h>
#include "fmgr.h"
#include "utils/geo_decls.h"

PG_MODULE_MAGIC;

/* by value */
PG_FUNCTION_INFO_V1(add_one);

Datum
add_one(PG_FUNCTION_ARGS)
{
    int32   arg = PG_GETARG_INT32(0);

    PG_RETURN_INT32(arg + 1);
}

/* by reference, fixed length */
PG_FUNCTION_INFO_V1(add_one_float8);

Datum
add_one_float8(PG_FUNCTION_ARGS)
{
    /* The macros for FLOAT8 hide its pass-by-reference nature. */
    float8   arg = PG_GETARG_FLOAT8(0);

    PG_RETURN_FLOAT8(arg + 1.0);
}

PG_FUNCTION_INFO_V1(makepoint);

Datum
makepoint(PG_FUNCTION_ARGS)
{
    /* Here, the pass-by-reference nature of Point is not hidden. */
    Point    *pointx = PG_GETARG_POINT_P(0);
    Point    *pointy = PG_GETARG_POINT_P(1);
    Point    *new_point = (Point *) palloc(sizeof(Point));

    new_point->x = pointx->x;
    new_point->y = pointy->y;

    PG_RETURN_POINT_P(new_point);
}

/* by reference, variable length */
Extending SQL

PG_FUNCTION_INFO_V1(copytext);

Datum

    copytext(PG_FUNCTION_ARGS)
    {
        text   *t = PG_GETARG_TEXT_PP(0);

        /*
        * VARSIZE_ANY_EXHDR is the size of the struct in bytes, minus
        * VARHDRSZ or VARHDRSZ_SHORT of its header. Construct the
        * copy with a
        * full-length header.
        */
        text   *new_t = (text *) palloc(VARSIZE_ANY_EXHDR(t) + VARHDRSZ);
        SET_VARSIZE(new_t, VARSIZE_ANY_EXHDR(t) + VARHDRSZ);

        /*
        * VARDATA is a pointer to the data region of the new struct.
        * The source
        * could be a short datum, so retrieve its data through
        * VARDATA_ANY.
        */
        memcpy((void *) VARDATA(new_t), /* destination */
               (void *) VARDATA_ANY(t), /* source */
               VARSIZE_ANY_EXHDR(t));   /* how many bytes */

        PG_RETURN_TEXT_P(new_t);
    }

PG_FUNCTION_INFO_V1(concat_text);

Datum

    concat_text(PG_FUNCTION_ARGS)
    {
        text   *arg1 = PG_GETARG_TEXT_PP(0);
        text   *arg2 = PG_GETARG_TEXT_PP(1);
        int32   arg1_size = VARSIZE_ANY_EXHDR(arg1);
        int32   arg2_size = VARSIZE_ANY_EXHDR(arg2);
        int32   new_text_size = arg1_size + arg2_size + VARHDRSZ;
        text    *new_text = (text *) palloc(new_text_size);

        SET_VARSIZE(new_text, new_text_size);
        memcpy(VARDATA(new_text), VARDATA_ANY(arg1), arg1_size);
        memcpy(VARDATA(new_text) + arg1_size, VARDATA_ANY(arg2),
               arg2_size);

        PG_RETURN_TEXT_P(new_text);
    }

Supposing that the above code has been prepared in file funcs.c and compiled into a shared object, we could define the functions to PostgreSQL with commands like this:

    CREATE FUNCTION add_one(integer) RETURNS integer
        AS 'DIRECTORY/funcs', 'add_one'
        LANGUAGE C STRICT;
-- note overloading of SQL function name "add_one"
CREATE FUNCTION add_one(double precision) RETURNS double precision
    AS 'DIRECTORY/funcs', 'add_one_float8'
    LANGUAGE C STRICT;

CREATE FUNCTION makepoint(point, point) RETURNS point
    AS 'DIRECTORY/funcs', 'makepoint'
    LANGUAGE C STRICT;

CREATE FUNCTION copytext(text) RETURNS text
    AS 'DIRECTORY/funcs', 'copytext'
    LANGUAGE C STRICT;

CREATE FUNCTION concat_text(text, text) RETURNS text
    AS 'DIRECTORY/funcs', 'concat_text'
    LANGUAGE C STRICT;

Here, DIRECTORY stands for the directory of the shared library file (for instance the PostgreSQL tutorial directory, which contains the code for the examples used in this section). (Better style would be to use just 'funcs' in the AS clause, after having added DIRECTORY to the search path. In any case, we can omit the system-specific extension for a shared library, commonly .so.)

Notice that we have specified the functions as "strict", meaning that the system should automatically assume a null result if any input value is null. By doing this, we avoid having to check for null inputs in the function code. Without this, we'd have to check for null values explicitly, using PG_ARGISNULL().

At first glance, the version-1 coding conventions might appear to be just pointless obscurantism, over using plain C calling conventions. They do however allow to deal with NULLable arguments/return values, and “toasted” (compressed or out-of-line) values.

The macro PG_ARGISNULL(n) allows a function to test whether each input is null. (Of course, doing this is only necessary in functions not declared “strict”.) As with the PG_GETARG_xxx() macros, the input arguments are counted beginning at zero. Note that one should refrain from executing PG_GETARG_xxx() until one has verified that the argument isn't null. To return a null result, execute PG_RETURN_NULL(); this works in both strict and nonstrict functions.

Other options provided by the version-1 interface are two variants of the PG_GETARG_xxx() macros. The first of these, PG_GETARG_xxx_COPY(), guarantees to return a copy of the specified argument that is safe for writing into. (The normal macros will sometimes return a pointer to a value that is physically stored in a table, which must not be written to. Using the PG_GETARG_xxx_COPY() macros guarantees a writable result.) The second variant consists of the PG_GETARG_xxx_SLICE() macros which take three arguments. The first is the number of the function argument (as above). The second and third are the offset and length of the segment to be returned. Offsets are counted from zero, and a negative length requests that the remainder of the value be returned. These macros provide more efficient access to parts of large values in the case where they have storage type “external”. (The storage type of a column can be specified using ALTER TABLE tablename ALTER COLUMN colname SET STORAGE storagetype. storagetype is one of plain, external, extended, or main.)

Finally, the version-1 function call conventions make it possible to return set results (Section 38.10.8) and implement trigger functions (Chapter 39) and procedural-language call handlers (Chapter 56). For more details see src/backend/utils/fmgr/README in the source distribution.

38.10.4. Writing Code

Before we turn to the more advanced topics, we should discuss some coding rules for PostgreSQL C-language functions. While it might be possible to load functions written in languages other than C into
Extending SQL

PostgreSQL, this is usually difficult (when it is possible at all) because other languages, such as C++, FORTRAN, or Pascal often do not follow the same calling convention as C. That is, other languages do not pass argument and return values between functions in the same way. For this reason, we will assume that your C-language functions are actually written in C.

The basic rules for writing and building C functions are as follows:

• Use `pg_config --includedir-server` to find out where the PostgreSQL server header files are installed on your system (or the system that your users will be running on).

• Compiling and linking your code so that it can be dynamically loaded into PostgreSQL always requires special flags. See Section 38.10.5 for a detailed explanation of how to do it for your particular operating system.

• Remember to define a “magic block” for your shared library, as described in Section 38.10.1.

• When allocating memory, use the PostgreSQL functions `palloc` and `pfree` instead of the corresponding C library functions `malloc` and `free`. The memory allocated by `palloc` will be freed automatically at the end of each transaction, preventing memory leaks.

• Always zero the bytes of your structures using `memset` (or allocate them with `palloc0` in the first place). Even if you assign to each field of your structure, there might be alignment padding (holes in the structure) that contain garbage values. Without this, it’s difficult to support hash indexes or hash joins, as you must pick out only the significant bits of your data structure to compute a hash. The planner also sometimes relies on comparing constants via bitwise equality, so you can get undesirable planning results if logically-equivalent values aren’t bitwise equal.

• Most of the internal PostgreSQL types are declared in `postgres.h`, while the function manager interfaces (`PG_FUNCTION_ARGS`, etc.) are in `fmgr.h`, so you will need to include at least these two files. For portability reasons it’s best to include `postgres.h` first, before any other system or user header files. Including `postgres.h` will also include `elog.h` and `palloc.h` for you.

• Symbol names defined within object files must not conflict with each other or with symbols defined in the PostgreSQL server executable. You will have to rename your functions or variables if you get error messages to this effect.

38.10.5. Compiling and Linking Dynamically-loaded Functions

Before you are able to use your PostgreSQL extension functions written in C, they must be compiled and linked in a special way to produce a file that can be dynamically loaded by the server. To be precise, a shared library needs to be created.

For information beyond what is contained in this section you should read the documentation of your operating system, in particular the manual pages for the C compiler, `cc`, and the link editor, `ld`. In addition, the PostgreSQL source code contains several working examples in the `contrib` directory. If you rely on these examples you will make your modules dependent on the availability of the PostgreSQL source code, however.

Creating shared libraries is generally analogous to linking executables: first the source files are compiled into object files, then the object files are linked together. The object files need to be created as position-independent code (PIC), which conceptually means that they can be placed at an arbitrary location in memory when they are loaded by the executable. (Object files intended for executables are usually not compiled that way.) The command to link a shared library contains special flags to distinguish it from linking an executable (at least in theory — on some systems the practice is much uglier).

In the following examples we assume that your source code is in a file `foo.c` and we will create a shared library `foo.so`. The intermediate object file will be called `foo.o` unless otherwise noted. A shared library can contain more than one object file, but we only use one here.
FreeBSD

The compiler flag to create PIC is \(-fPIC\). To create shared libraries the compiler flag is \(-\text{shared}\).

\[
gcc -fPIC -c foo.c
gcc -shared -o foo.so foo.o
\]

This is applicable as of version 3.0 of FreeBSD.

HP-UX

The compiler flag of the system compiler to create PIC is \(+z\). When using GCC it's \(-fPIC\). The linker flag for shared libraries is \(-b\). So:

\[
cc +z -c foo.c
\]
or:

\[
gcc -fPIC -c foo.c
\]

and then:

\[
ld -b -o foo.sl foo.o
\]

HP-UX uses the extension \(.sl\) for shared libraries, unlike most other systems.

Linux

The compiler flag to create PIC is \(-fPIC\). The compiler flag to create a shared library is \(-\text{shared}\). A complete example looks like this:

\[
cc -fPIC -c foo.c
c -shared -o foo.so foo.o
\]

macOS

Here is an example. It assumes the developer tools are installed.

\[
cc -c foo.c
c -bundle -flat_namespace -undefined suppress -o foo.so foo.o
\]

NetBSD

The compiler flag to create PIC is \(-fPIC\). For ELF systems, the compiler with the flag \(-\text{shared}\) is used to link shared libraries. On the older non-ELF systems, \(ld -B\text{shareable}\) is used.

\[
gcc -fPIC -c foo.c
gcc -shared -o foo.so foo.o
\]

OpenBSD

The compiler flag to create PIC is \(-fPIC\). \(ld -B\text{shareable}\) is used to link shared libraries.

\[
gcc -fPIC -c foo.c
ld -B\text{shareable} -o foo.so foo.o
\]
The compiler flag to create PIC is `-KPIC` with the Sun compiler and `-fPIC` with GCC. To link shared libraries, the compiler option is `-G` with either compiler or alternatively `-shared` with GCC.

```bash
c -KPIC -c foo.c
c -G -o foo.so foo.o
```
or

```bash
gcc -fPIC -c foo.c
gcc -G -o foo.so foo.o
```

**Tip**

If this is too complicated for you, you should consider using GNU Libtool\(^1\), which hides the platform differences behind a uniform interface.

The resulting shared library file can then be loaded into PostgreSQL. When specifying the file name to the `CREATE FUNCTION` command, one must give it the name of the shared library file, not the intermediate object file. Note that the system’s standard shared-library extension (usually `.so` or `.sl`) can be omitted from the `CREATE FUNCTION` command, and normally should be omitted for best portability.

Refer back to Section 38.10.1 about where the server expects to find the shared library files.

### 38.10.6. Composite-type Arguments

Composite types do not have a fixed layout like C structures. Instances of a composite type can contain null fields. In addition, composite types that are part of an inheritance hierarchy can have different fields than other members of the same inheritance hierarchy. Therefore, PostgreSQL provides a function interface for accessing fields of composite types from C.

Suppose we want to write a function to answer the query:

```sql
SELECT name, c_overpaid(emp, 1500) AS overpaid
FROM emp
WHERE name = 'Bill' OR name = 'Sam';
```

Using the version-1 calling conventions, we can define `c_overpaid` as:

```c
#include "postgres.h"
#include "executor/execparser.h" /* for GetAttributeByName() */

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(c_overpaid);

Datum
c_overpaid(PG_FUNCTION_ARGS)
{
```

\(^1\)http://www.gnu.org/software/libtool/
HeapTupleHeader  t = PG_GETARG_HEAPTUPLEHEADER(0);
int32            limit = PG_GETARG_INT32(1);
bool isnull;
Datum salary;

salary = GetAttributeByName(t, "salary", &isnull);
if (isnull)
    PG_RETURN_BOOL(false);
/* Alternatively, we might prefer to do PG_RETURN_NULL() for
null salary. */

    PG_RETURN_BOOL(DatumGetInt32(salary) > limit);
}

GetAttributeByName is the PostgreSQL system function that returns attributes out of the specified row. It has three arguments: the argument of type HeapTupleHeader passed into the function, the name of the desired attribute, and a return parameter that tells whether the attribute is null. GetAttributeByName returns a Datum value that you can convert to the proper data type by using the appropriate DatumGetXXX() macro. Note that the return value is meaningless if the null flag is set; always check the null flag before trying to do anything with the result.

There is also GetAttributeByNum, which selects the target attribute by column number instead of name.

The following command declares the function c_overpaid in SQL:

```
CREATE FUNCTION c_overpaid(emp, integer) RETURNS boolean
    AS 'DIRECTORY/funcs', 'c_overpaid'
    LANGUAGE C STRICT;
```

Notice we have used STRICT so that we did not have to check whether the input arguments were NULL.

### 38.10.7. Returning Rows (Composite Types)

To return a row or composite-type value from a C-language function, you can use a special API that provides macros and functions to hide most of the complexity of building composite data types. To use this API, the source file must include:

```
#include "funcapi.h"
```

There are two ways you can build a composite data value (henceforth a “tuple”): you can build it from an array of Datum values, or from an array of C strings that can be passed to the input conversion functions of the tuple's column data types. In either case, you first need to obtain or construct a TupleDesc descriptor for the tuple structure. When working with Datums, you pass the TupleDesc to BlessTupleDesc, and then call heap_form_tuple for each row. When working with C strings, you pass the TupleDesc to TupleDescGetAttInMetadata, and then call BuildTupleFromCStrings for each row. In the case of a function returning a set of tuples, the setup steps can all be done once during the first call of the function.

Several helper functions are available for setting up the needed TupleDesc. The recommended way to do this in most functions returning composite values is to call:

```
TypeFuncClass get_call_result_type(FunctionCallInfo fcinfo,
    Oid *resultTypeId,
    TupleDesc *resultTupleDesc)
```
passing the same fcinfo struct passed to the calling function itself. (This of course requires that you use the version-1 calling conventions.) resultTypeId can be specified as NULL or as the address of a local variable to receive the function's result type OID. resultTupleDesc should be the address of a local TupleDesc variable. Check that the result is TYPEFUNC_COMPOSITE; if so, resultTupleDesc has been filled with the needed TupleDesc. (If it is not, you can report an error along the lines of "function returning record called in context that cannot accept type record").

### Tip
get_call_result_type can resolve the actual type of a polymorphic function result; so it is useful in functions that return scalar polymorphic results, not only functions that return composites. The resultTypeId output is primarily useful for functions returning polymorphic scalars.

### Note
get_call_result_type has a sibling get_expr_result_type, which can be used to resolve the expected output type for a function call represented by an expression tree. This can be used when trying to determine the result type from outside the function itself. There is also get_func_result_type, which can be used when only the function's OID is available. However these functions are not able to deal with functions declared to return record, and get_func_result_type cannot resolve polymorphic types, so you should preferentially use get_call_result_type.

Older, now-deprecated functions for obtaining TupleDescs are:

TupleDesc RelationNameGetTupleDesc(const char *relname)
to get a TupleDesc for the row type of a named relation, and:

TupleDesc TypeGetTupleDesc(Oid typeoid, List *colaliases)
to get a TupleDesc based on a type OID. This can be used to get a TupleDesc for a base or composite type. It will not work for a function that returns record, however, and it cannot resolve polymorphic types.

Once you have a TupleDesc, call:

TupleDesc BlessTupleDesc(TupleDesc tupdesc)
if you plan to work with Datums, or:

AttInMetadata *TupleDescGetAttInMetadata(TupleDesc tupdesc)
if you plan to work with C strings. If you are writing a function returning set, you can save the results of these functions in the FuncCallContext structure — use the tuple_desc or attinmeta field respectively.

When working with Datums, use:

HeapTuple heap_form_tuple(TupleDesc tupdesc, Datum *values, bool *isnull)
to build a HeapTuple given user data in Datum form.

When working with C strings, use:

```c
HeapTuple BuildTupleFromCStrings(AttInMetadata *attinmeta, char **values)
```
to build a HeapTuple given user data in C string form. `values` is an array of C strings, one for each attribute of the return row. Each C string should be in the form expected by the input function of the attribute data type. In order to return a null value for one of the attributes, the corresponding pointer in the `values` array should be set to `NULL`. This function will need to be called again for each row you return.

Once you have built a tuple to return from your function, it must be converted into a Datum. Use:

```c
HeapTupleGetDatum(HeapTuple tuple)
```
to convert a HeapTuple into a valid Datum. This Datum can be returned directly if you intend to return just a single row, or it can be used as the current return value in a set-returning function.

An example appears in the next section.

### 38.10.8. Returning Sets

There is also a special API that provides support for returning sets (multiple rows) from a C-language function. A set-returning function must follow the version-1 calling conventions. Also, source files must include `funcapi.h`, as above.

A set-returning function (SRF) is called once for each item it returns. The SRF must therefore save enough state to remember what it was doing and return the next item on each call. The structure `FuncCallContext` is provided to help control this process. Within a function, `fcinfo->flinfo->fn_extra` is used to hold a pointer to `FuncCallContext` across calls.

```c
typedef struct FuncCallContext
{
    /*
    * Number of times we've been called before
    *
    * call_cntr is initialized to 0 for you by
    * SRF_FIRSTCALL_INIT(), and
    * incremented for you every time SRF_RETURN_NEXT() is called.
    */
    uint64 call_cntr;

    /*
    * OPTIONAL maximum number of calls
    *
    * max_calls is here for convenience only and setting it is optional.
    * If not set, you must provide alternative means to know when the
    * function is done.
    */
    uint64 max_calls;

    /*
    * OPTIONAL pointer to result slot
    *
* This is obsolete and only present for backward compatibility, viz,
  * user-defined SRFs that use the deprecated
  TupleDescGetSlot().
  */
  TupleTableSlot *slot;

  /*
   * OPTIONAL pointer to miscellaneous user-provided context
   * information
   * user_fctx is for use as a pointer to your own data to retain
   * arbitrary context information between calls of your
   * function.
   */
  void *user_fctx;

  /*
   * OPTIONAL pointer to struct containing attribute type input
   * metadata
   * attinmeta is for use when returning tuples (i.e., composite
   * data types)
   * and is not used when returning base data types. It is only
   * needed
   * if you intend to use BuildTupleFromCStrings() to create the
   * return
   * tuple.
   */
  AttInMetadata *attinmeta;

  /*
   * memory context used for structures that must live for
   * multiple calls
   * multi_call_memory_ctx is set by SRF_FIRSTCALL_INIT() for
   * you, and used
   * by SRF_RETURN_DONE() for cleanup. It is the most appropriate
   * memory
   * context for any memory that is to be reused across multiple
   * calls
   * of the SRF.
   */
  MemoryContext multi_call_memory_ctx;

  /*
   * OPTIONAL pointer to struct containing tuple description
   * tuple_desc is for use when returning tuples (i.e., composite
   * data types)
   * and is only needed if you are going to build the tuples with
   * heap_form_tuple() rather than with BuildTupleFromCStrings().
   * Note that
   * the TupleDesc pointer stored here should usually have been
   * run through
   * BlessTupleDesc() first.
   */
  TupleDesc tuple_desc;
An SRF uses several functions and macros that automatically manipulate the `FuncCallContext` structure (and expect to find it via `fn_extra`). Use:

```c
SRF_IS_FIRSTCALL()
```

to determine if your function is being called for the first or a subsequent time. On the first call (only) use:

```c
SRF_FIRSTCALL_INIT()
```
to initialize the `FuncCallContext`. On every function call, including the first, use:

```c
SRF_PERCALL_SETUP()
```
to properly set up for using the `FuncCallContext` and clearing any previously returned data left over from the previous pass.

If your function has data to return, use:

```c
SRF_RETURN_NEXT(funcctx, result)
```
to return it to the caller. (`result` must be of type `Datum`, either a single value or a tuple prepared as described above.) Finally, when your function is finished returning data, use:

```c
SRF_RETURN_DONE(funcctx)
```
to clean up and end the SRF.

The memory context that is current when the SRF is called is a transient context that will be cleared between calls. This means that you do not need to call `pfree` on everything you allocated using `palloc`; it will go away anyway. However, if you want to allocate any data structures to live across calls, you need to put them somewhere else. The memory context referenced by `multi_call_memory_ctx` is a suitable location for any data that needs to survive until the SRF is finished running. In most cases, this means that you should switch into `multi_call_memory_ctx` while doing the first-call setup.

**Warning**

While the actual arguments to the function remain unchanged between calls, if you detoast the argument values (which is normally done transparently by the `PG_GETARG_xxx` macro) in the transient context then the detoasted copies will be freed on each cycle. Accordingly, if you keep references to such values in your `user_fctx`, you must either copy them into the `multi_call_memory_ctx` after detoasting, or ensure that you detoast the values only in that context.

A complete pseudo-code example looks like the following:

```c
Datum
my_set_returning_function(PG_FUNCTION_ARGS)
{
```
Extending SQL

FuncCallContext *funcctx;
Datum result;

/* further declarations as needed */

if (SRF_IS_FIRSTCALL())
{
    MemoryContext oldcontext;

    funcctx = SRF_FIRSTCALL_INIT();
    oldcontext = MemoryContextSwitchTo(funcctx->multi_call_memory_ctx);
    /* One-time setup code appears here: */
    user code
    if returning composite
        build TupleDesc, and perhaps AttInMetadata
    endif returning composite
    user code
    MemoryContextSwitchTo(oldcontext);
}

/* Each-time setup code appears here: */
user code
funcctx = SRF_PERCALL_SETUP();
user code

/* this is just one way we might test whether we are done: */
if (funcctx->call_cntr < funcctx->max_calls)
{
    /* Here we want to return another item: */
    user code
    obtain result Datum
    SRF_RETURN_NEXT(funcctx, result);
}
else
{
    /* Here we are done returning items and just need to clean up: */
    user code
    SRF_RETURN_DONE(funcctx);
}

A complete example of a simple SRF returning a composite type looks like:

PG_FUNCTION_INFO_V1(retcomposite);

Datum retcomposite(PG_FUNCTION_ARGS)
{
    FuncCallContext *funcctx;
    int call_cntr;
    int max_calls;
    TupleDesc tupdesc;
    AttInMetadata *attinmeta;

    /* stuff done only on the first call of the function */
    if (SRF_IS_FIRSTCALL())
    {
MemoryContext oldcontext;

/* create a function context for cross-call persistence */
funcctx = SRF_FIRSTCALL_INIT();

/* switch to memory context appropriate for multiple function calls */
oldcontext = MemoryContextSwitchTo(funcctx->multi_call_memory_ctx);

/* total number of tuples to be returned */
funcctx->max_calls = PG_GETARG_UINT32(0);

/* Build a tuple descriptor for our result type */
if (get_call_result_type(fcinfo, NULL, &tupdesc) != TYPEFUNC_COMPOSITE)
    ereport(ERROR,
            (errcode(ERRCODE_FEATURE_NOT_SUPPORTED),
             errmsg("function returning record called in context "
                    "that cannot accept type record")));

/* generate attribute metadata needed later to produce tuples from raw */
    * C strings */
attinmeta = TupleDescGetAttInMetadata(tupdesc);
    funcctx->attinmeta = attinmeta;

    MemoryContextSwitchTo(oldcontext);
}

/* stuff done on every call of the function */
funcctx = SRF_PERCALL_SETUP();

call_cntr = funcctx->call_cntr;
max_calls = funcctx->max_calls;
attinmeta = funcctx->attinmeta;

if (call_cntr < max_calls) /* do when there is more left to send */
{
    char **values;
    HeapTuple tuple;
    Datum result;

    /* Prepare a values array for building the returned tuple. 
     * This should be an array of C strings which will 
     * be processed later by the type input functions. */
    values = (char **) palloc(3 * sizeof(char *));
    values[0] = (char *) palloc(16 * sizeof(char));
    values[1] = (char *) palloc(16 * sizeof(char));
    values[2] = (char *) palloc(16 * sizeof(char));

    /* do something with values */
One way to declare this function in SQL is:

```sql
CREATE TYPE __retcomposite AS (f1 integer, f2 integer, f3 integer);

CREATE OR REPLACE FUNCTION retcomposite(integer, integer)
    RETURNS SETOF __retcomposite
    AS 'filename', 'retcomposite'
    LANGUAGE C IMMUTABLE STRICT;
```

A different way is to use OUT parameters:

```sql
CREATE OR REPLACE FUNCTION retcomposite(IN integer, IN integer,
    OUT f1 integer, OUT f2 integer, OUT f3 integer)
    RETURNS SETOF record
    AS 'filename', 'retcomposite'
    LANGUAGE C IMMUTABLE STRICT;
```

Notice that in this method the output type of the function is formally an anonymous record type.

The directory `contrib/tablefunc` module in the source distribution contains more examples of set-returning functions.

### 38.10.9. Polymorphic Arguments and Return Types

C-language functions can be declared to accept and return the polymorphic types `anyelement`, `anyarray`, `anynonarray`, `anyenum`, and `anyrange`. See Section 38.2.5 for a more detailed explanation of polymorphic functions. When function arguments or return types are defined as polymorphic types, the function author cannot know in advance what data type it will be called with, or need to return. There are two routines provided in `fmgr.h` to allow a version-1 C function to discover the actual data types of its arguments and the type it is expected to return. The routines are called `get_fn_expr_rettype(FmgrInfo *flinfo)` and `get_fn_expr_argtype(FmgrIn-
fo *flinfo, int argnum). They return the result or argument type OID, or InvalidOid if
the information is not available. The structure flinfo is normally accessed as fcinfo->flinfo.
The parameter argnum is zero based. get_call_result_type can also be used as an alterna-
tive to get_fn_expr_rettype. There is also get_fn_expr_variadic, which can be used
to find out whether variadic arguments have been merged into an array. This is primarily useful for
VARIADIC "any" functions, since such merging will always have occurred for variadic functions
taking ordinary array types.

For example, suppose we want to write a function to accept a single element of any type, and return
a one-dimensional array of that type:

PG_FUNCTION_INFO_V1(make_array);
Datum
make_array(PG_FUNCTION_ARGS)
{
   ArrayType  *result;
   Oid         element_type = get_fn_expr_argtype(fcinfo->flinfo, 0);
   Datum       element;
   bool        isnull;
   int16       typlen;
   bool        typbyval;
   char        typalign;
   int         ndims;
   int         dims[MAXDIM];
   int         lbs[MAXDIM];

   if (!OidIsValid(element_type))
      elog(ERROR, "could not determine data type of input");

   /* get the provided element, being careful in case it's NULL */
   isnull = PG_ARGISNULL(0);
   if (isnull)
      element = (Datum) 0;
   else
      element = PG_GETARG_DATUM(0);

   /* we have one dimension */
   ndims = 1;
   /* and one element */
   dims[0] = 1;
   /* and lower bound is 1 */
   lbs[0] = 1;

   /* get required info about the element type */
   get_typlenbyvalalign(element_type, &typlen, &typbyval, &typalign);

   /* now build the array */
   result = construct_md_array(&element, &isnull, ndims, dims, lbs,
                                element_type, typlen, typbyval, typalign);

   PG_RETURN_ARRAYTYPE_P(result);
}

The following command declares the function make_array in SQL:
CREATE FUNCTION make_array(anyelement) RETURNS anyarray
AS 'DIRECTORY/funcs', 'make_array'
LANGUAGE C IMMUTABLE;

There is a variant of polymorphism that is only available to C-language functions: they can be declared
to take parameters of type "any". (Note that this type name must be double-quoted, since it's also a
SQL reserved word.) This works like anyelement except that it does not constrain different "any"
arguments to be the same type, nor do they help determine the function's result type. A C-language
function can also declare its final parameter to be VARIADIC "any". This will match one or more
actual arguments of any type (not necessarily the same type). These arguments will not be gathered into
an array as happens with normal variadic functions; they will just be passed to the function separately.
The PG_NARGS () macro and the methods described above must be used to determine the number of
actual arguments and their types when using this feature. Also, users of such a function might wish
to use the VARIADIC keyword in their function call, with the expectation that the function would
treat the array elements as separate arguments. The function itself must implement that behavior if
wanted, after using get_fn_expr_variadic to detect that the actual argument was marked with
VARIADIC.

38.10.10. Transform Functions

Some function calls can be simplified during planning based on properties specific to the function.
For example, int4mul(n, 1) could be simplified to just n. To define such function-specific opti-
mizations, write a transform function and place its OID in the protransform field of the primary
function's pg_proc entry. The transform function must have the SQL signature protransfor-
m(internal) RETURNS internal. The argument, actually FuncExpr *, is a dummy node
representing a call to the primary function. If the transform function's study of the expression tree
proves that a simplified expression tree can substitute for all possible concrete calls represented there-
by, build and return that simplified expression. Otherwise, return a NULL pointer (not a SQL null).

We make no guarantee that PostgreSQL will never call the primary function in cases that the transform
function could simplify. Ensure rigorous equivalence between the simplified expression and an actual
call to the primary function.

Currently, this facility is not exposed to users at the SQL level because of security concerns, so it is
only practical to use for optimizing built-in functions.

38.10.11. Shared Memory and LWLocks

Add-ins can reserve LWLocks and an allocation of shared memory on server startup. The add-in's
shared library must be preloaded by specifying it in shared_preload_libraries. Shared memory is re-
served by calling:

void RequestAddinShmemSpace(int size)

from your _PG_init function.

LWLocks are reserved by calling:

void RequestNamedLWLockTranche(const char *tranche_name, int
num_lwlocks)

from _PG_init. This will ensure that an array of num_lwlocks LWLocks is available under the
name tranche_name. Use GetNamedLWLockTranche to get a pointer to this array.

To avoid possible race-conditions, each backend should use the LWLock AddinShmemInitLock
when connecting to and initializing its allocation of shared memory, as shown here:
Extending SQL

static mystruct *ptr = NULL;

if (!ptr)
{
  bool found;

  LWLockAcquire(AddinShmemInitLock, LW_EXCLUSIVE);
  ptr = ShmemInitStruct("my struct name", size, &found);
  if (!found)
  {
    initialize contents of shmem area;
    acquire any requested LWLocks using:
    ptr->locks = GetNamedLWLockTranche("my tranche name");
  }
  LWLockRelease(AddinShmemInitLock);
}

38.10.12. Using C++ for Extensibility

Although the PostgreSQL backend is written in C, it is possible to write extensions in C++ if these guidelines are followed:

• All functions accessed by the backend must present a C interface to the backend; these C functions can then call C++ functions. For example, extern C linkage is required for backend-accessed functions. This is also necessary for any functions that are passed as pointers between the backend and C++ code.

• Free memory using the appropriate deallocation method. For example, most backend memory is allocated using palloc(), so use pfree() to free it. Using C++ delete in such cases will fail.

• Prevent exceptions from propagating into the C code (use a catch-all block at the top level of all extern C functions). This is necessary even if the C++ code does not explicitly throw any exceptions, because events like out-of-memory can still throw exceptions. Any exceptions must be caught and appropriate errors passed back to the C interface. If possible, compile C++ with -fno-exceptions to eliminate exceptions entirely; in such cases, you must check for failures in your C++ code, e.g. check for NULL returned by new().

• If calling backend functions from C++ code, be sure that the C++ call stack contains only plain old data structures (POD). This is necessary because backend errors generate a distant longjmp() that does not properly unroll a C++ call stack with non-POD objects.

In summary, it is best to place C++ code behind a wall of extern C functions that interface to the backend, and avoid exception, memory, and call stack leakage.

38.11. User-defined Aggregates

Aggregate functions in PostgreSQL are defined in terms of state values and state transition functions. That is, an aggregate operates using a state value that is updated as each successive input row is processed. To define a new aggregate function, one selects a data type for the state value, an initial value for the state, and a state transition function. The state transition function takes the previous state value and the aggregate's input value(s) for the current row, and returns a new state value. A final function can also be specified, in case the desired result of the aggregate is different from the data that needs to be kept in the running state value. The final function takes the ending state value and returns whatever is wanted as the aggregate result. In principle, the transition and final functions are just ordinary functions that could also be used outside the context of the aggregate. (In practice, it's
often helpful for performance reasons to create specialized transition functions that can only work when called as part of an aggregate.

Thus, in addition to the argument and result data types seen by a user of the aggregate, there is an internal state-value data type that might be different from both the argument and result types.

If we define an aggregate that does not use a final function, we have an aggregate that computes a running function of the column values from each row. sum is an example of this kind of aggregate. sum starts at zero and always adds the current row's value to its running total. For example, if we want to make a sum aggregate to work on a data type for complex numbers, we only need the addition function for that data type. The aggregate definition would be:

```sql
CREATE AGGREGATE sum (complex)
(
    sfunc = complex_add,
    stype = complex,
    initcond = '(0,0)'
);
```

which we might use like this:

```sql
SELECT sum(a) FROM test_complex;
```

<table>
<thead>
<tr>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>(34,53.9)</td>
</tr>
</tbody>
</table>

(Notice that we are relying on function overloading: there is more than one aggregate named sum, but PostgreSQL can figure out which kind of sum applies to a column of type complex.)

The above definition of sum will return zero (the initial state value) if there are no nonnull input values. Perhaps we want to return null in that case instead — the SQL standard expects sum to behave that way. We can do this simply by omitting the initcond phrase, so that the initial state value is null. Ordinarily this would mean that the sfunc would need to check for a null state-value input. But for sum and some other simple aggregates like max and min, it is sufficient to insert the first nonnull input value into the state variable and then start applying the transition function at the second nonnull input value. PostgreSQL will do that automatically if the initial state value is null and the transition function is marked “strict” (i.e., not to be called for null inputs).

Another bit of default behavior for a “strict” transition function is that the previous state value is retained unchanged whenever a null input value is encountered. Thus, null values are ignored. If you need some other behavior for null inputs, do not declare your transition function as strict; instead code it to test for null inputs and do whatever is needed.

avg (average) is a more complex example of an aggregate. It requires two pieces of running state: the sum of the inputs and the count of the number of inputs. The final result is obtained by dividing these quantities. Average is typically implemented by using an array as the state value. For example, the built-in implementation of avg(float8) looks like:

```sql
CREATE AGGREGATE avg (float8)
(
    sfunc = float8_accum,
    stype = float8[],
    finalfunc = float8_avg,
    initcond = '{0,0,0}'
);
```
Extending SQL

Note

**float8_accum** requires a three-element array, not just two elements, because it accumulates the sum of squares as well as the sum and count of the inputs. This is so that it can be used for some other aggregates as well as **avg**.

Aggregate function calls in SQL allow **DISTINCT** and **ORDER BY** options that control which rows are fed to the aggregate's transition function and in what order. These options are implemented behind the scenes and are not the concern of the aggregate's support functions.

For further details see the CREATE AGGREGATE command.

### 38.11.1. Moving-Aggregate Mode

Aggregate functions can optionally support **moving-aggregate mode**, which allows substantially faster execution of aggregate functions within windows with moving frame starting points. (See Section 3.5 and Section 4.2.8 for information about use of aggregate functions as window functions.) The basic idea is that in addition to a normal “forward” transition function, the aggregate provides an **inverse transition function**, which allows rows to be removed from the aggregate's running state value when they exit the window frame. For example a **sum** aggregate, which uses addition as the forward transition function, would use subtraction as the inverse transition function. Without an inverse transition function, the window function mechanism must recalculate the aggregate from scratch each time the frame starting point moves, resulting in run time proportional to the number of input rows times the average frame length. With an inverse transition function, the run time is only proportional to the number of input rows.

The inverse transition function is passed the current state value and the aggregate input value(s) for the earliest row included in the current state. It must reconstruct what the state value would have been if the given input row had never been aggregated, but only the rows following it. This sometimes requires that the forward transition function keep more state than is needed for plain aggregation mode. Therefore, the moving-aggregate mode uses a completely separate implementation from the plain mode: it has its own state data type, its own forward transition function, and its own final function if needed. These can be the same as the plain mode's data type and functions, if there is no need for extra state.

As an example, we could extend the **sum** aggregate given above to support moving-aggregate mode like this:

```sql
CREATE AGGREGATE sum (complex) (
    sfunc = complex_add,
   stype = complex,
    initcond = '(0,0)',
    msfunc = complex_add,
    minvfunc = complex_sub,
    mstype = complex,
    minitcond = '(0,0)' );
```

The parameters whose names begin with **m** define the moving-aggregate implementation. Except for the inverse transition function **minvfunc**, they correspond to the plain-aggregate parameters without **m**.

The forward transition function for moving-aggregate mode is not allowed to return null as the new state value. If the inverse transition function returns null, this is taken as an indication that the inverse
function cannot reverse the state calculation for this particular input, and so the aggregate calculation
will be redone from scratch for the current frame starting position. This convention allows moving-ag-
ggregate mode to be used in situations where there are some infrequent cases that are impractical to
reverse out of the running state value. The inverse transition function can “punt” on these cases, and
yet still come out ahead so long as it can work for most cases. As an example, an aggregate working
with floating-point numbers might choose to punt when a NaN (not a number) input has to be removed
from the running state value.

When writing moving-aggregate support functions, it is important to be sure that the inverse transition
function can reconstruct the correct state value exactly. Otherwise there might be user-visible differ-
ences in results depending on whether the moving-aggregate mode is used. An example of an aggregate
for which adding an inverse transition function seems easy at first, yet where this requirement cannot
be met is sum over float4 or float8 inputs. A naive declaration of \texttt{sum(float8)} could be

\begin{verbatim}
CREATE AGGREGATE unsafe_sum (float8)
{
   stype = float8,
   sfunc = float8pl,
   mstype = float8,
   msfunc = float8pl,
   minvfunc = float8mi
};
\end{verbatim}

This aggregate, however, can give wildly different results than it would have without the inverse
transition function. For example, consider

\begin{verbatim}
SELECT unsafe_sum(x) OVER (ORDER BY n ROWS BETWEEN CURRENT ROW AND 1
FOLLOWING)
FROM (VALUES (1, 1.0e20::float8),
(2, 1.0::float8)) AS v (n,x);
\end{verbatim}

This query returns 0 as its second result, rather than the expected answer of 1. The cause is the limited
precision of floating-point values: adding 1 to 1e20 results in 1e20 again, and so subtracting 1e20
from that yields 0, not 1. Note that this is a limitation of floating-point arithmetic in general, not a
limitation of PostgreSQL.

38.11.2. Polymorphic and Variadic Aggregates

Aggregate functions can use polymorphic state transition functions or final functions, so that the same
functions can be used to implement multiple aggregates. See Section 38.2.5 for an explanation of
polymorphic functions. Going a step further, the aggregate function itself can be specified with poly-
morphic input type(s) and state type, allowing a single aggregate definition to serve for multiple input
data types. Here is an example of a polymorphic aggregate:

\begin{verbatim}
CREATE AGGREGATE array_accum (anyelement)
{
   sfunc = array_append,
   stype = anyarray,
   initcond = '{}'
};
\end{verbatim}

Here, the actual state type for any given aggregate call is the array type having the actual input type as
elements. The behavior of the aggregate is to concatenate all the inputs into an array of that type. (Note:
the built-in aggregate \texttt{array_agg} provides similar functionality, with better performance than this
definition would have.)
Here's the output using two different actual data types as arguments:

```
SELECT attrelid::regclass, array_accum(attname)
FROM pg_attribute
WHERE attnum > 0 AND attrelid = 'pg_tablespace'::regclass
GROUP BY attrelid;
```

<table>
<thead>
<tr>
<th>attrelid</th>
<th>array_accum</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_tablespace</td>
<td>{spcname, spcowner, spcacl, spcoptions}</td>
</tr>
</tbody>
</table>

(1 row)

```
SELECT attrelid::regclass, array_accum(atttypid::regtype)
FROM pg_attribute
WHERE attnum > 0 AND attrelid = 'pg_tablespace'::regclass
GROUP BY attrelid;
```

<table>
<thead>
<tr>
<th>attrelid</th>
<th>array_accum</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_tablespace</td>
<td>{name, oid, aclitem[], text[]}</td>
</tr>
</tbody>
</table>

(1 row)

Ordinarily, an aggregate function with a polymorphic result type has a polymorphic state type, as in the above example. This is necessary because otherwise the final function cannot be declared sensibly: it would need to have a polymorphic result type but no polymorphic argument type, which CREATE FUNCTION will reject on the grounds that the result type cannot be deduced from a call. But sometimes it is inconvenient to use a polymorphic state type. The most common case is where the aggregate support functions are to be written in C and the state type should be declared as internal because there is no SQL-level equivalent for it. To address this case, it is possible to declare the final function as taking extra “dummy” arguments that match the input arguments of the aggregate. Such dummy arguments are always passed as null values since no specific value is available when the final function is called. Their only use is to allow a polymorphic final function's result type to be connected to the aggregate's input type(s). For example, the definition of the built-in aggregate `array_agg` is equivalent to

```
CREATE FUNCTION array_agg_transfn(internal, anynonarray)
RETURNS internal ...;
CREATE FUNCTION array_agg_finalfn(internal, anynonarray)
RETURNS anyarray ...;

CREATE AGGREGATE array_agg (anynonarray)
{
    sfunc = array_agg_transfn,
    stype = internal,
    finalfunc = array_agg_finalfn,
    finalfunc_extra
};
```

Here, the `finalfunc_extra` option specifies that the final function receives, in addition to the state value, extra dummy argument(s) corresponding to the aggregate's input argument(s). The extra `anynonarray` argument allows the declaration of `array_agg_finalfn` to be valid.

An aggregate function can be made to accept a varying number of arguments by declaring its last argument as a VARIADIC array, in much the same fashion as for regular functions; see Section 38.5.5. The aggregate's transition function(s) must have the same array type as their last argument. The transition function(s) typically would also be marked VARIADIC, but this is not strictly required.
Variadic aggregates are easily misused in connection with the ORDER BY option (see Section 4.2.7), since the parser cannot tell whether the wrong number of actual arguments have been given in such a combination. Keep in mind that everything to the right of ORDER BY is a sort key, not an argument to the aggregate. For example, in

```
SELECT myaggregate(a ORDER BY a, b, c) FROM ...
```

the parser will see this as a single aggregate function argument and three sort keys. However, the user might have intended

```
SELECT myaggregate(a, b, c ORDER BY a) FROM ...
```

If `myaggregate` is variadic, both these calls could be perfectly valid.

For the same reason, it's wise to think twice before creating aggregate functions with the same names and different numbers of regular arguments.

### 38.11.3. Ordered-Set Aggregates

The aggregates we have been describing so far are “normal” aggregates. PostgreSQL also supports *ordered-set aggregates*, which differ from normal aggregates in two key ways. First, in addition to ordinary aggregated arguments that are evaluated once per input row, an ordered-set aggregate can have “direct” arguments that are evaluated only once per aggregation operation. Second, the syntax for the ordinary aggregated arguments specifies a sort ordering for them explicitly. An ordered-set aggregate is usually used to implement a computation that depends on a specific row ordering, for instance rank or percentile, so that the sort ordering is a required aspect of any call. For example, the built-in definition of `percentile_disc` is equivalent to:

```sql
CREATE FUNCTION ordered_set_transition(internal, anyelement)
    RETURNS internal ...;
CREATE FUNCTION percentile_disc_final(internal, float8, anyelement)
    RETURNS anyelement ...;

CREATE AGGREGATE percentile_disc (float8 ORDER BY anyelement)
{
    sfunc = ordered_set_transition,
    stype = internal,
    finalfunc = percentile_disc_final,
    finalfunc_extra
};
```

This aggregate takes a `float8` direct argument (the percentile fraction) and an aggregated input that can be of any sortable data type. It could be used to obtain a median household income like this:

```
SELECT percentile_disc(0.5) WITHIN GROUP (ORDER BY income) FROM households;
```

```
percentile_disc
-----------------
50489
```

Here, 0.5 is a direct argument; it would make no sense for the percentile fraction to be a value varying across rows.
Unlike the case for normal aggregates, the sorting of input rows for an ordered-set aggregate is not done behind the scenes, but is the responsibility of the aggregate's support functions. The typical implementation approach is to keep a reference to a "tuplesort" object in the aggregate's state value, feed the incoming rows into that object, and then complete the sorting and read out the data in the final function. This design allows the final function to perform special operations such as injecting additional "hypothetical" rows into the data to be sorted. While normal aggregates can often be implemented with support functions written in PL/pgSQL or another PL language, ordered-set aggregates generally have to be written in C, since their state values aren't definable as any SQL data type. (In the above example, notice that the state value is declared as type internal — this is typical.)

Also, because the final function performs the sort, it is not possible to continue adding input rows by executing the transition function again later. This means the final function is not READ_ONLY; it must be declared in CREATE AGGREGATE as READ_WRITE, or as SHAREABLE if it's possible for additional final-function calls to make use of the already-sorted state.

The state transition function for an ordered-set aggregate receives the current state value plus the aggregated input values for each row, and returns the updated state value. This is the same definition as for normal aggregates, but note that the direct arguments (if any) are not provided. The final function receives the last state value, the values of the direct arguments if any, and (if finalfunc_extra is specified) null values corresponding to the aggregated input(s). As with normal aggregates, finalfunc_extra is only really useful if the aggregate is polymorphic; then the extra dummy argument(s) are needed to connect the final function's result type to the aggregate's input type(s).

Currently, ordered-set aggregates cannot be used as window functions, and therefore there is no need for them to support moving-aggregate mode.

### 38.11.4. Partial Aggregation

Optionally, an aggregate function can support partial aggregation. The idea of partial aggregation is to run the aggregate's state transition function over different subsets of the input data independently, and then to combine the state values resulting from those subsets to produce the same state value that would have resulted from scanning all the input in a single operation. This mode can be used for parallel aggregation by having different worker processes scan different portions of a table. Each worker produces a partial state value, and at the end those state values are combined to produce a final state value. (In the future this mode might also be used for purposes such as combining aggregations over local and remote tables; but that is not implemented yet.)

To support partial aggregation, the aggregate definition must provide a combine function, which takes two values of the aggregate's state type (representing the results of aggregating over two subsets of the input rows) and produces a new value of the state type, representing what the state would have been after aggregating over the combination of those sets of rows. It is unspecified what the relative order of the input rows from the two sets would have been. This means that it's usually impossible to define a useful combine function for aggregates that are sensitive to input row order.

As simple examples, MAX and MIN aggregates can be made to support partial aggregation by specifying the combine function as the same greater-of-two or lesser-of-two comparison function that is used as their transition function. SUM aggregates just need an addition function as combine function. (Again, this is the same as their transition function, unless the state value is wider than the input data type.)

The combine function is treated much like a transition function that happens to take a value of the state type, not of the underlying input type, as its second argument. In particular, the rules for dealing with null values and strict functions are similar. Also, if the aggregate definition specifies a non-null initcond, keep in mind that that will be used not only as the initial state for each partial aggregation run, but also as the initial state for the combine function, which will be called to combine each partial result into that state.

If the aggregate's state type is declared as internal, it is the combine function's responsibility that its result is allocated in the correct memory context for aggregate state values. This means in particular that when the first input is NULL it's invalid to simply return the second input, as that value will be in the wrong context and will not have sufficient lifespan.
When the aggregate's state type is declared as internal, it is usually also appropriate for the aggregate definition to provide a serialization function and a deserialization function, which allow such a state value to be copied from one process to another. Without these functions, parallel aggregation cannot be performed, and future applications such as local/remote aggregation will probably not work either.

A serialization function must take a single argument of type internal and return a result of type bytea, which represents the state value packaged up into a flat blob of bytes. Conversely, a deserialization function reverses that conversion. It must take two arguments of types bytea and internal, and return a result of type internal. (The second argument is unused and is always zero, but it is required for type-safety reasons.) The result of the deserialization function should simply be allocated in the current memory context, as unlike the combine function's result, it is not long-lived.

Worth noting also is that for an aggregate to be executed in parallel, the aggregate itself must be marked PARALLEL SAFE. The parallel-safety markings on its support functions are not consulted.

### 38.11.5. Support Functions for Aggregates

A function written in C can detect that it is being called as an aggregate support function by calling `AggCheckCallContext`, for example:

```c
if (AggCheckCallContext(fcinfo, NULL))
```

One reason for checking this is that when it is true, the first input must be a temporary state value and can therefore safely be modified in-place rather than allocating a new copy. See `int8inc()` for an example. (While aggregate transition functions are always allowed to modify the transition value in-place, aggregate final functions are generally discouraged from doing so; if they do so, the behavior must be declared when creating the aggregate. See CREATE AGGREGATE for more detail.)

The second argument of `AggCheckCallContext` can be used to retrieve the memory context in which aggregate state values are being kept. This is useful for transition functions that wish to use “expanded” objects (see Section 38.12.1) as their state values. On first call, the transition function should return an expanded object whose memory context is a child of the aggregate state context, and then keep returning the same expanded object on subsequent calls. See `array_append()` for an example. (`array_append()` is not the transition function of any built-in aggregate, but it is written to behave efficiently when used as transition function of a custom aggregate.)

Another support routine available to aggregate functions written in C is `AggGetAggref`, which returns the `Aggref` parse node that defines the aggregate call. This is mainly useful for ordered-set aggregates, which can inspect the substructure of the `Aggref` node to find out what sort ordering they are supposed to implement. Examples can be found in `orderedsetaggs.c` in the PostgreSQL source code.

### 38.12. User-defined Types

As described in Section 38.2, PostgreSQL can be extended to support new data types. This section describes how to define new base types, which are data types defined below the level of the SQL language. Creating a new base type requires implementing functions to operate on the type in a low-level language, usually C.

The examples in this section can be found in `complex.sql` and `complex.c` in the `src/tutorial` directory of the source distribution. See the README file in that directory for instructions about running the examples.

A user-defined type must always have input and output functions. These functions determine how the type appears in strings (for input by the user and output to the user) and how the type is organized in memory. The input function takes a null-terminated character string as its argument and returns the internal (in memory) representation of the type. The output function takes the internal representation
of the type as argument and returns a null-terminated character string. If we want to do anything more with the type than merely store it, we must provide additional functions to implement whatever operations we’d like to have for the type.

Suppose we want to define a type `complex` that represents complex numbers. A natural way to represent a complex number in memory would be the following C structure:

```c
typedef struct Complex {
    double x;
    double y;
} Complex;
```

We will need to make this a pass-by-reference type, since it’s too large to fit into a single `Datum` value.

As the external string representation of the type, we choose a string of the form `(x, y)`.

The input and output functions are usually not hard to write, especially the output function. But when defining the external string representation of the type, remember that you must eventually write a complete and robust parser for that representation as your input function. For instance:

```c
PG_FUNCTION_INFO_V1(complex_in);
Datum complex_in(PG_FUNCTION_ARGS)
{
    char       *str = PG_GETARG_CSTRING(0);
    double      x,
                 y;
    Complex    *result;

    if (sscanf(str, " ( %lf , %lf )", &x, &y) != 2)
        ereport(ERROR,
                (errcode(ERRCODE_INVALID_TEXT_REPRESENTATION),
                errmsg("invalid input syntax for complex: "%s\",
                        str));

    result = (Complex *) palloc(sizeof(Complex));
    result->x = x;
    result->y = y;
    PG_RETURN_POINTER(result);
}
```

The output function can simply be:

```c
PG_FUNCTION_INFO_V1(complex_out);
Datum complex_out(PG_FUNCTION_ARGS)
{
    Complex   *complex = (Complex *) PG_GETARG_POINTER(0);
    char       *result;

    result = psprintf("(%g,%g)", complex->x, complex->y);
    PG_RETURN_CSTRING(result);
}
```
You should be careful to make the input and output functions inverses of each other. If you do not, you will have severe problems when you need to dump your data into a file and then read it back in. This is a particularly common problem when floating-point numbers are involved.

Optionally, a user-defined type can provide binary input and output routines. Binary I/O is normally faster but less portable than textual I/O. As with textual I/O, it is up to you to define exactly what the external binary representation is. Most of the built-in data types try to provide a machine-independent binary representation. For complex, we will piggy-back on the binary I/O converters for type \texttt{float8}:

```c
PG_FUNCTION_INFO_V1(complex_recv);

Datum
complex_recv(PG_FUNCTION_ARGS)
{
    StringInfo  buf = (StringInfo) PG_GETARG_POINTER(0);
    Complex    *result;

    result = (Complex *) malloc(sizeof(Complex));
    result->x = pq_getmsgfloat8(buf);
    result->y = pq_getmsgfloat8(buf);
    PG_RETURN_POINTER(result);
}

PG_FUNCTION_INFO_V1(complex_send);

Datum
complex_send(PG_FUNCTION_ARGS)
{
    Complex    *complex = (Complex *) PG_GETARG_POINTER(0);
    StringInfoData buf;

    pq_begintypsend(&buf);
    pq_sendfloat8(&buf, complex->x);
    pq_sendfloat8(&buf, complex->y);
    PG_RETURN_BYTEA_P(pq_endtypsend(&buf));
}
```

Once we have written the I/O functions and compiled them into a shared library, we can define the complex type in SQL. First we declare it as a shell type:

```sql
CREATE TYPE complex;
```

This serves as a placeholder that allows us to reference the type while defining its I/O functions. Now we can define the I/O functions:

```sql
CREATE FUNCTION complex_in(cstring)
    RETURNS complex
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;

CREATE FUNCTION complex_out(complex)
    RETURNS cstring
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;
```
CREATE FUNCTION complex_recv(internal)
    RETURNS complex
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;

CREATE FUNCTION complex_send(complex)
    RETURNS bytea
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;

Finally, we can provide the full definition of the data type:

CREATE TYPE complex (  
    internallength = 16,
    input = complex_in,
    output = complex_out,
    receive = complex_recv,
    send = complex_send,
    alignment = double
);

When you define a new base type, PostgreSQL automatically provides support for arrays of that type. The array type typically has the same name as the base type with the underscore character (_) prepended.

Once the data type exists, we can declare additional functions to provide useful operations on the data type. Operators can then be defined atop the functions, and if needed, operator classes can be created to support indexing of the data type. These additional layers are discussed in following sections.

If the internal representation of the data type is variable-length, the internal representation must follow the standard layout for variable-length data: the first four bytes must be a char[4] field which is never accessed directly (customarily named vl_len_). You must use the SET_VARSIZE() macro to store the total size of the datum (including the length field itself) in this field and VARSIZE() to retrieve it. (These macros exist because the length field may be encoded depending on platform.)

For further details see the description of the CREATE TYPE command.

**38.12.1. TOAST Considerations**

If the values of your data type vary in size (in internal form), it's usually desirable to make the data type TOAST-able (see Section 68.2). You should do this even if the values are always too small to be compressed or stored externally, because TOAST can save space on small data too, by reducing header overhead.

To support TOAST storage, the C functions operating on the data type must always be careful to unpack any toasted values they are handed by using PG_DETOAST_DATUM. (This detail is customarily hidden by defining type-specific GETARG_DATATYPE_P macros.) Then, when running the CREATE TYPE command, specify the internal length as variable and select some appropriate storage option other than plain.

If data alignment is unimportant (either just for a specific function or because the data type specifies byte alignment anyway) then it's possible to avoid some of the overhead of PG_DETOAST_DATUM. You can use PG_DETOAST_DATUM_PACKED instead (customarily hidden by defining a GETARG_DATATYPE_PP macro) and using the macros VARSIZE_ANY_EXHDR and VAR_DATA_ANY to access a potentially-packed datum. Again, the data returned by these macros is not aligned even if the data type definition specifies an alignment. If the alignment is important you must go through the regular PG_DETOAST_DATUM interface.
Extending SQL

Note

Older code frequently declares `vl_len_` as an `int32` field instead of `char[4]`. This is OK as long as the struct definition has other fields that have at least `int32` alignment. But it is dangerous to use such a struct definition when working with a potentially unaligned datum; the compiler may take it as license to assume the datum actually is aligned, leading to core dumps on architectures that are strict about alignment.

Another feature that’s enabled by TOAST support is the possibility of having an expanded in-memory data representation that is more convenient to work with than the format that is stored on disk. The regular or “flat” varlena storage format is ultimately just a blob of bytes; it cannot for example contain pointers, since it may get copied to other locations in memory. For complex data types, the flat format may be quite expensive to work with, so PostgreSQL provides a way to “expand” the flat format into a representation that is more suited to computation, and then pass that format in-memory between functions of the data type.

To use expanded storage, a data type must define an expanded format that follows the rules given in `src/include/utils/expandeddatum.h`, and provide functions to “expand” a flat varlena value into expanded format and “flatten” the expanded format back to the regular varlena representation. Then ensure that all C functions for the data type can accept either representation, possibly by converting one into the other immediately upon receipt. This does not require fixing all existing functions for the data type at once, because the standard `PG_DETOAST_DATUM` macro is defined to convert expanded inputs into regular flat format. Therefore, existing functions that work with the flat varlena format will continue to work, though slightly inefficiently, with expanded inputs; they need not be converted until and unless better performance is important.

C functions that know how to work with an expanded representation typically fall into two categories: those that can only handle expanded format, and those that can handle either expanded or flat varlena inputs. The former are easier to write but may be less efficient overall, because converting a flat input to expanded form for use by a single function may cost more than is saved by operating on the expanded format. When only expanded format need be handled, conversion of flat inputs to expanded form can be hidden inside an argument-fetching macro, so that the function appears no more complex than one working with traditional varlena input. To handle both types of input, write an argument-fetching function that will detoast external, short-header, and compressed varlena inputs, but not expanded inputs. Such a function can be defined as returning a pointer to a union of the flat varlena format and the expanded format. Callers can use the `VARATT_IS_EXPANDED_HEADER()` macro to determine which format they received.

The TOAST infrastructure not only allows regular varlena values to be distinguished from expanded values, but also distinguishes “read-write” and “read-only” pointers to expanded values. C functions that only need to examine an expanded value, or will only change it in safe and non-semantically-visible ways, need not care which type of pointer they receive. C functions that produce a modified version of an input value are allowed to modify an expanded input value in-place if they receive a read-write pointer, but must not modify the input if they receive a read-only pointer; in that case they have to copy the value first, producing a new value to modify. A C function that has constructed a new expanded value should always return a read-write pointer to it. Also, a C function that is modifying a read-write expanded value in-place should take care to leave the value in a sane state if it fails partway through.

For examples of working with expanded values, see the standard array infrastructure, particularly `src/backend/utils/adt/array_expanded.c`.

### 38.13. User-defined Operators

Every operator is “syntactic sugar” for a call to an underlying function that does the real work; so you must first create the underlying function before you can create the operator. However, an operator is not merely syntactic sugar, because it carries additional information that helps the query planner
optimize queries that use the operator. The next section will be devoted to explaining that additional
information.

PostgreSQL supports left unary, right unary, and binary operators. Operators can be overloaded; that
is, the same operator name can be used for different operators that have different numbers and types
of operands. When a query is executed, the system determines the operator to call from the number
and types of the provided operands.

Here is an example of creating an operator for adding two complex numbers. We assume we’ve already
created the definition of type complex (see Section 38.12). First we need a function that does the
work, then we can define the operator:

```sql
CREATE FUNCTION complex_add(complex, complex)
  RETURNS complex
  AS 'filename', 'complex_add'
  LANGUAGE C IMMUTABLE STRICT;
CREATE OPERATOR + (
  leftarg = complex,
  rightarg = complex,
  function = complex_add,
  commutator = +
);
```

Now we could execute a query like this:

```sql
SELECT (a + b) AS c FROM test_complex;
```

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>5.2,6.05</td>
<td>(133.42,144.95)</td>
</tr>
</tbody>
</table>

We’ve shown how to create a binary operator here. To create unary operators, just omit one of left-
arg (for left unary) or rightarg (for right unary). The function clause and the argument clauses
are the only required items in CREATE OPERATOR. The commutator clause shown in the example
is an optional hint to the query optimizer. Further details about commutator and other optimizer
hints appear in the next section.

### 38.14. Operator Optimization Information

A PostgreSQL operator definition can include several optional clauses that tell the system useful things
about how the operator behaves. These clauses should be provided whenever appropriate, because they
can make for considerable speedups in execution of queries that use the operator. But if you provide
them, you must be sure that they are right! Incorrect use of an optimization clause can result in slow
queries, subtly wrong output, or other Bad Things. You can always leave out an optimization clause
if you are not sure about it; the only consequence is that queries might run slower than they need to.

Additional optimization clauses might be added in future versions of PostgreSQL. The ones described
here are all the ones that release 11.1 understands.

#### 38.14.1. commutator

The commutator clause, if provided, names an operator that is the commutator of the operator being
defined. We say that operator A is the commutator of operator B if (x A y) equals (y B x) for all
possible input values x, y. Notice that B is also the commutator of A. For example, operators < and >
for a particular data type are usually each others' commutators, and operator + is usually commutative with itself. But operator – is usually not commutative with anything.

The left operand type of a commutable operator is the same as the right operand type of its commutator, and vice versa. So the name of the commutator operator is all that PostgreSQL needs to be given to look up the commutator, and that's all that needs to be provided in the COMMUTATOR clause.

It's critical to provide commutator information for operators that will be used in indexes and join clauses, because this allows the query optimizer to “flip around” such a clause to the forms needed for different plan types. For example, consider a query with a WHERE clause like tab1.x = tab2.y, where tab1.x and tab2.y are of a user-defined type, and suppose that tab2.y is indexed. The optimizer cannot generate an index scan unless it can determine how to flip the clause around to tab2.y = tab1.x, because the index-scan machinery expects to see the indexed column on the left of the operator it is given. PostgreSQL will not simply assume that this is a valid transformation — the creator of the = operator must specify that it is valid, by marking the operator with commutator information.

When you are defining a self-commutative operator, you just do it. When you are defining a pair of commutative operators, things are a little trickier: how can the first one to be defined refer to the other one, which you haven't defined yet? There are two solutions to this problem:

- One way is to omit the COMMUTATOR clause in the first operator that you define, and then provide one in the second operator's definition. Since PostgreSQL knows that commutative operators come in pairs, when it sees the second definition it will automatically go back and fill in the missing COMMUTATOR clause in the first definition.

- The other, more straightforward way is just to include COMMUTATOR clauses in both definitions. When PostgreSQL processes the first definition and realizes that COMMUTATOR refers to a nonexistent operator, the system will make a dummy entry for that operator in the system catalog. This dummy entry will have valid data only for the operator name, left and right operand types, and result type, since that's all that PostgreSQL can deduce at this point. The first operator's catalog entry will link to this dummy entry. Later, when you define the second operator, the system updates the dummy entry with the additional information from the second definition. If you try to use the dummy operator before it's been filled in, you'll just get an error message.

### 38.14.2. NEGATOR

The NEGATOR clause, if provided, names an operator that is the negator of the operator being defined. We say that operator A is the negator of operator B if both return Boolean results and (x A y) equals NOT (x B y) for all possible inputs x, y. Notice that B is also the negator of A. For example, < and >= are a negator pair for most data types. An operator can never validly be its own negator.

Unlike commutators, a pair of unary operators could validly be marked as each other's negators; that would mean (A x) equals NOT (B x) for all x, or the equivalent for right unary operators.

An operator's negator must have the same left and/or right operand types as the operator to be defined, so just as with COMMUTATOR, only the operator name need be given in the NEGATOR clause.

Providing a negator is very helpful to the query optimizer since it allows expressions like NOT (x = y) to be simplified into x <> y. This comes up more often than you might think, because NOT operations can be inserted as a consequence of other rearrangements.

Pairs of negator operators can be defined using the same methods explained above for commutator pairs.

### 38.14.3. RESTRICT

The RESTRICT clause, if provided, names a restriction selectivity estimation function for the operator. (Note that this is a function name, not an operator name.) RESTRICT clauses only make sense for
binary operators that return boolean. The idea behind a restriction selectivity estimator is to guess what fraction of the rows in a table will satisfy a WHERE-clause condition of the form:

column OP constant

for the current operator and a particular constant value. This assists the optimizer by giving it some idea of how many rows will be eliminated by WHERE clauses that have this form. (What happens if the constant is on the left, you might be wondering? Well, that's one of the things that COMMUTATOR is for...)

Writing new restriction selectivity estimation functions is far beyond the scope of this chapter, but fortunately you can usually just use one of the system's standard estimators for many of your own operators. These are the standard restriction estimators:

eqsel for =
neqsel for <>
scalarltsel for <
scalarsel for <=
scalargtssel for >
scalargesel for >=

You can frequently get away with using either eqsel or neqsel for operators that have very high or very low selectivity, even if they aren't really equality or inequality. For example, the approximate-equality geometric operators use eqsel on the assumption that they'll usually only match a small fraction of the entries in a table.

You can use scalarltssel, scalarsel, scalargtssel and scalargesel for comparisons on data types that have some sensible means of being converted into numeric scalars for range comparisons. If possible, add the data type to those understood by the function convert_to_scalar() in src/backend/utils/adt/selfuncks.c. (Eventually, this function should be replaced by per-data-type functions identified through a column of the pg_type system catalog; but that hasn't happened yet.) If you do not do this, things will still work, but the optimizer's estimates won't be as good as they could be.

There are additional selectivity estimation functions designed for geometric operators in src/backend/utils/adt/geo_selfuncks.c: areasel, positionsel, and contsel. At this writing these are just stubs, but you might want to use them (or even better, improve them) anyway.

38.14.4. JOIN

The JOIN clause, if provided, names a join selectivity estimation function for the operator. (Note that this is a function name, not an operator name.) JOIN clauses only make sense for binary operators that return boolean. The idea behind a join selectivity estimator is to guess what fraction of the rows in a pair of tables will satisfy a WHERE-clause condition of the form:

table1.column1 OP table2.column2

for the current operator. As with the RESTRICT clause, this helps the optimizer very substantially by letting it figure out which of several possible join sequences is likely to take the least work.

As before, this chapter will make no attempt to explain how to write a join selectivity estimator function, but will just suggest that you use one of the standard estimators if one is applicable:
eqjoinsel for =
neqjoinsel for <>
scalarltjoinsel for <
scalarlejoinsel for <=
scalargtjoinsel for >
scalargejoinsel for >=
areajoin for 2D area-based comparisons
positionjoin for 2D position-based comparisons
contjoin for 2D containment-based comparisons

38.14.5. HASHES

The HASHES clause, if present, tells the system that it is permissible to use the hash join method for a join based on this operator. HASHES only makes sense for a binary operator that returns boolean, and in practice the operator must represent equality for some data type or pair of data types.

The assumption underlying hash join is that the join operator can only return true for pairs of left and right values that hash to the same hash code. If two values get put in different hash buckets, the join will never compare them at all, implicitly assuming that the result of the join operator must be false. So it never makes sense to specify HASHES for operators that do not represent some form of equality. In most cases it is only practical to support hashing for operators that take the same data type on both sides. However, sometimes it is possible to design compatible hash functions for two or more data types; that is, functions that will generate the same hash codes for “equal” values, even though the values have different representations. For example, it’s fairly simple to arrange this property when hashing integers of different widths.

To be marked HASHES, the join operator must appear in a hash index operator family. This is not enforced when you create the operator, since of course the referencing operator family couldn’t exist yet. But attempts to use the operator in hash joins will fail at run time if no such operator family exists. The system needs the operator family to find the data-type-specific hash function(s) for the operator’s input data type(s). Of course, you must also create suitable hash functions before you can create the operator family.

Care should be exercised when preparing a hash function, because there are machine-dependent ways in which it might fail to do the right thing. For example, if your data type is a structure in which there might be uninteresting pad bits, you cannot simply pass the whole structure to hash_any. (Unless you write your other operators and functions to ensure that the unused bits are always zero, which is the recommended strategy.) Another example is that on machines that meet the IEEE floating-point standard, negative zero and positive zero are different values (different bit patterns) but they are defined to compare equal. If a float value might contain negative zero then extra steps are needed to ensure it generates the same hash value as positive zero.

A hash-joinable operator must have a commutator (itself if the two operand data types are the same, or a related equality operator if they are different) that appears in the same operator family. If this is not the case, planner errors might occur when the operator is used. Also, it is a good idea (but not strictly required) for a hash operator family that supports multiple data types to provide equality operators for every combination of the data types; this allows better optimization.

Note

The function underlying a hash-joinable operator must be marked immutable or stable. If it is volatile, the system will never attempt to use the operator for a hash join.

Note

If a hash-joinable operator has an underlying function that is marked strict, the function must also be complete: that is, it should return true or false, never null, for any two nonnull inputs. If this rule is not followed, hash-optimization of IN operations might generate wrong results. (Specifically, IN might return false where the correct answer according to the standard would be null; or it might yield an error complaining that it wasn’t prepared for a null result.)
38.14.6. MERGES

The MERGES clause, if present, tells the system that it is permissible to use the merge-join method for a join based on this operator. MERGES only makes sense for a binary operator that returns boolean, and in practice the operator must represent equality for some data type or pair of data types.

Merge join is based on the idea of sorting the left- and right-hand tables into order and then scanning them in parallel. So, both data types must be capable of being fully ordered, and the join operator must be one that can only succeed for pairs of values that fall at the “same place” in the sort order. In practice this means that the join operator must behave like equality. But it is possible to merge-join two distinct data types so long as they are logically compatible. For example, the smallint-versus-integer equality operator is merge-joinable. We only need sorting operators that will bring both data types into a logically compatible sequence.

To be marked MERGES, the join operator must appear as an equality member of a btree index operator family. This is not enforced when you create the operator, since of course the referencing operator family couldn't exist yet. But the operator will not actually be used for merge joins unless a matching operator family can be found. The MERGES flag thus acts as a hint to the planner that it's worth looking for a matching operator family.

A merge-joinable operator must have a commutator (itself if the two operand data types are the same, or a related equality operator if they are different) that appears in the same operator family. If this is not the case, planner errors might occur when the operator is used. Also, it is a good idea (but not strictly required) for a btree operator family that supports multiple data types to provide equality operators for every combination of the data types; this allows better optimization.

Note

The function underlying a merge-joinable operator must be marked immutable or stable. If it is volatile, the system will never attempt to use the operator for a merge join.

38.15. Interfacing Extensions To Indexes

The procedures described thus far let you define new types, new functions, and new operators. However, we cannot yet define an index on a column of a new data type. To do this, we must define an operator class for the new data type. Later in this section, we will illustrate this concept in an example: a new operator class for the B-tree index method that stores and sorts complex numbers in ascending absolute value order.

Operator classes can be grouped into operator families to show the relationships between semantically compatible classes. When only a single data type is involved, an operator class is sufficient, so we'll focus on that case first and then return to operator families.

38.15.1. Index Methods and Operator Classes

The pg_am table contains one row for every index method (internally known as access method). Support for regular access to tables is built into PostgreSQL, but all index methods are described in pg_am. It is possible to add a new index access method by writing the necessary code and then creating an entry in pg_am — but that is beyond the scope of this chapter (see Chapter 61).

The routines for an index method do not directly know anything about the data types that the index method will operate on. Instead, an operator class identifies the set of operations that the index method needs to use to work with a particular data type. Operator classes are so called because one thing they specify is the set of WHERE-clause operators that can be used with an index (i.e., can be converted into an index-scan qualification). An operator class can also specify some support function that are
needed by the internal operations of the index method, but do not directly correspond to any WHERE-clause operator that can be used with the index.

It is possible to define multiple operator classes for the same data type and index method. By doing this, multiple sets of indexing semantics can be defined for a single data type. For example, a B-tree index requires a sort ordering to be defined for each data type it works on. It might be useful for a complex-number data type to have one B-tree operator class that sorts the data by complex absolute value, another that sorts by real part, and so on. Typically, one of the operator classes will be deemed most commonly useful and will be marked as the default operator class for that data type and index method.

The same operator class name can be used for several different index methods (for example, both B-tree and hash index methods have operator classes named \texttt{int4_ops}), but each such class is an independent entity and must be defined separately.

### 38.15.2. Index Method Strategies

The operators associated with an operator class are identified by “strategy numbers”, which serve to identify the semantics of each operator within the context of its operator class. For example, B-trees impose a strict ordering on keys, lesser to greater, and so operators like “less than” and “greater than or equal to” are interesting with respect to a B-tree. Because PostgreSQL allows the user to define operators, PostgreSQL cannot look at the name of an operator (e.g., \texttt{<} or \texttt{>=}) and tell what kind of comparison it is. Instead, the index method defines a set of “strategies”, which can be thought of as generalized operators. Each operator class specifies which actual operator corresponds to each strategy for a particular data type and interpretation of the index semantics.

The B-tree index method defines five strategies, shown in Table 38.2.

#### Table 38.2. B-tree Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>less than</td>
<td>1</td>
</tr>
<tr>
<td>less than or equal</td>
<td>2</td>
</tr>
<tr>
<td>equal</td>
<td>3</td>
</tr>
<tr>
<td>greater than or equal</td>
<td>4</td>
</tr>
<tr>
<td>greater than</td>
<td>5</td>
</tr>
</tbody>
</table>

Hash indexes support only equality comparisons, and so they use only one strategy, shown in Table 38.3.

#### Table 38.3. Hash Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>equal</td>
<td>1</td>
</tr>
</tbody>
</table>

GiST indexes are more flexible: they do not have a fixed set of strategies at all. Instead, the “consistency” support routine of each particular GiST operator class interprets the strategy numbers however it likes. As an example, several of the built-in GiST index operator classes index two-dimensional geometric objects, providing the “R-tree” strategies shown in Table 38.4. Four of these are true two-dimensional tests (overlaps, same, contains, contained by); four of them consider only the X direction; and the other four provide the same tests in the Y direction.

#### Table 38.4. GiST Two-Dimensional “R-tree” Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>strictly left of</td>
<td>1</td>
</tr>
</tbody>
</table>
Extending SQL

SP-GiST indexes are similar to GiST indexes in flexibility: they don't have a fixed set of strategies. Instead the support routines of each operator class interpret the strategy numbers according to the operator class's definition. As an example, the strategy numbers used by the built-in operator classes for points are shown in Table 38.5.

Table 38.5. SP-GiST Point Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>does not extend to right of</td>
<td>2</td>
</tr>
<tr>
<td>overlaps</td>
<td>3</td>
</tr>
<tr>
<td>does not extend to left of</td>
<td>4</td>
</tr>
<tr>
<td>strictly right of</td>
<td>5</td>
</tr>
<tr>
<td>same</td>
<td>6</td>
</tr>
<tr>
<td>contains</td>
<td>7</td>
</tr>
<tr>
<td>contained by</td>
<td>8</td>
</tr>
<tr>
<td>does not extend above</td>
<td>9</td>
</tr>
<tr>
<td>strictly below</td>
<td>10</td>
</tr>
<tr>
<td>strictly above</td>
<td>11</td>
</tr>
<tr>
<td>does not extend below</td>
<td>12</td>
</tr>
</tbody>
</table>

GIN indexes are similar to GiST and SP-GiST indexes, in that they don't have a fixed set of strategies either. Instead the support routines of each operator class interpret the strategy numbers according to the operator class's definition. As an example, the strategy numbers used by the built-in operator class for arrays are shown in Table 38.6.

Table 38.6. GIN Array Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>overlap</td>
<td>1</td>
</tr>
<tr>
<td>contains</td>
<td>2</td>
</tr>
<tr>
<td>is contained by</td>
<td>3</td>
</tr>
<tr>
<td>equal</td>
<td>4</td>
</tr>
</tbody>
</table>

BRIN indexes are similar to GiST, SP-GiST and GIN indexes in that they don't have a fixed set of strategies either. Instead the support routines of each operator class interpret the strategy numbers according to the operator class's definition. As an example, the strategy numbers used by the built-in Minmax operator classes are shown in Table 38.7.

Table 38.7. BRIN Minmax Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>less than</td>
<td>1</td>
</tr>
</tbody>
</table>
Extending SQL

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>less than or equal</td>
<td>2</td>
</tr>
<tr>
<td>equal</td>
<td>3</td>
</tr>
<tr>
<td>greater than or equal</td>
<td>4</td>
</tr>
<tr>
<td>greater than</td>
<td>5</td>
</tr>
</tbody>
</table>

Notice that all the operators listed above return Boolean values. In practice, all operators defined as index method search operators must return type boolean, since they must appear at the top level of a WHERE clause to be used with an index. (Some index access methods also support ordering operators, which typically don't return Boolean values; that feature is discussed in Section 38.15.7.)

### 38.15.3. Index Method Support Routines

Strategies aren't usually enough information for the system to figure out how to use an index. In practice, the index methods require additional support routines in order to work. For example, the B-tree index method must be able to compare two keys and determine whether one is greater than, equal to, or less than the other. Similarly, the hash index method must be able to compute hash codes for key values. These operations do not correspond to operators used in qualifications in SQL commands; they are administrative routines used by the index methods, internally.

Just as with strategies, the operator class identifies which specific functions should play each of these roles for a given data type and semantic interpretation. The index method defines the set of functions it needs, and the operator class identifies the correct functions to use by assigning them to the “support function numbers” specified by the index method.

B-trees require a comparison support function, and allow two additional support functions to be supplied at the operator class author's option, as shown in Table 38.8. The requirements for these support functions are explained further in Section 63.3.

#### Table 38.8. B-tree Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Compare two keys and return an integer less than zero, zero, or greater than zero, indicating whether the first key is less than, equal to, or greater than the second</td>
<td>1</td>
</tr>
<tr>
<td>Return the addresses of C-callable sort support function(s) (optional)</td>
<td>2</td>
</tr>
<tr>
<td>Compare a test value to a base value plus/minus an offset, and return true or false according to the comparison result (optional)</td>
<td>3</td>
</tr>
</tbody>
</table>

Hash indexes require one support function, and allow a second one to be supplied at the operator class author's option, as shown in Table 38.9.

#### Table 38.9. Hash Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Compute the 32-bit hash value for a key</td>
<td>1</td>
</tr>
<tr>
<td>Compute the 64-bit hash value for a key given a 64-bit salt; if the salt is 0, the low 32 bits of the result must match the value that would have been computed by function 1 (optional)</td>
<td>2</td>
</tr>
</tbody>
</table>

GiST indexes have nine support functions, two of which are optional, as shown in Table 38.10. (For more information see Chapter 64.)
Table 38.10. GiST Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>1</td>
</tr>
<tr>
<td>union</td>
<td>compute union of a set of keys</td>
<td>2</td>
</tr>
<tr>
<td>compress</td>
<td>compute a compressed representation of a key or value to be indexed</td>
<td>3</td>
</tr>
<tr>
<td>decompress</td>
<td>compute a decompressed representation of a compressed key</td>
<td>4</td>
</tr>
<tr>
<td>penalty</td>
<td>compute penalty for inserting new key into subtree with given subtree's key</td>
<td>5</td>
</tr>
<tr>
<td>picksplit</td>
<td>determine which entries of a page are to be moved to the new page and compute the union keys for resulting pages</td>
<td>6</td>
</tr>
<tr>
<td>equal</td>
<td>compare two keys and return true if they are equal</td>
<td>7</td>
</tr>
<tr>
<td>distance</td>
<td>determine distance from key to query value (optional)</td>
<td>8</td>
</tr>
<tr>
<td>fetch</td>
<td>compute original representation of a compressed key for index-only scans (optional)</td>
<td>9</td>
</tr>
</tbody>
</table>

SP-GiST indexes require five support functions, as shown in Table 38.11. (For more information see Chapter 65.)

Table 38.11. SP-GiST Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>config</td>
<td>provide basic information about the operator class</td>
<td>1</td>
</tr>
<tr>
<td>choose</td>
<td>determine how to insert a new value into an inner tuple</td>
<td>2</td>
</tr>
<tr>
<td>picksplit</td>
<td>determine how to partition a set of values</td>
<td>3</td>
</tr>
<tr>
<td>inner_consistent</td>
<td>determine which sub-partitions need to be searched for a query</td>
<td>4</td>
</tr>
<tr>
<td>leaf_consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>5</td>
</tr>
</tbody>
</table>

GIN indexes have six support functions, three of which are optional, as shown in Table 38.12. (For more information see Chapter 66.)

Table 38.12. GIN Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>compare</td>
<td>compare two keys and return an integer less than zero, zero, or greater than zero, indicating whether the first key is less than,</td>
<td>1</td>
</tr>
</tbody>
</table>
BRIN indexes have four basic support functions, as shown in Table 38.13; those basic functions may require additional support functions to be provided. (For more information see Section 67.3.)

Table 38.13. BRIN Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>extractValue</td>
<td>extract keys from a value to be indexed</td>
<td>2</td>
</tr>
<tr>
<td>extractQuery</td>
<td>extract keys from a query condition</td>
<td>3</td>
</tr>
<tr>
<td>consistent</td>
<td>determine whether value matches query condition (Boolean variant) (optional if support function 6 is present)</td>
<td>4</td>
</tr>
<tr>
<td>comparePartial</td>
<td>compare partial key from query and key from index, and return an integer less than zero, zero, or greater than zero, indicating whether GIN should ignore this index entry, treat the entry as a match, or stop the index scan (optional)</td>
<td>5</td>
</tr>
<tr>
<td>triConsistent</td>
<td>determine whether value matches query condition (ternary variant) (optional if support function 4 is present)</td>
<td>6</td>
</tr>
</tbody>
</table>

Unlike search operators, support functions return whichever data type the particular index method expects; for example in the case of the comparison function for B-trees, a signed integer. The number and types of the arguments to each support function are likewise dependent on the index method. For B-tree and hash the comparison and hashing support functions take the same input data types as do the operators included in the operator class, but this is not the case for most GiST, SP-GiST, GIN, and BRIN support functions.

### 38.15.4. An Example

Now that we have seen the ideas, here is the promised example of creating a new operator class. (You can find a working copy of this example in `src/tutorial/complex.c` and `src/tutorial/complex.sql` in the source distribution.) The operator class encapsulates operators that sort complex numbers in absolute value order, so we choose the name `complex_abs_ops`. First, we need a set of operators. The procedure for defining operators was discussed in Section 38.13. For an operator class on B-trees, the operators we require are:
• absolute-value less-than (strategy 1)
• absolute-value less-than-or-equal (strategy 2)
• absolute-value equal (strategy 3)
• absolute-value greater-than-or-equal (strategy 4)
• absolute-value greater-than (strategy 5)

The least error-prone way to define a related set of comparison operators is to write the B-tree comparison support function first, and then write the other functions as one-line wrappers around the support function. This reduces the odds of getting inconsistent results for corner cases. Following this approach, we first write:

```c
#define Mag(c)  ((c)->x*(c)->x + (c)->y*(c)->y)
static int complex_abs_cmp_internal(Complex *a, Complex *b)
{
    double amag = Mag(a),
            bmag = Mag(b);

    if (amag < bmag)
        return -1;
    if (amag > bmag)
        return 1;
    return 0;
}
```

Now the less-than function looks like:

```c
PG_FUNCTION_INFO_V1(complex_abs_lt);
Datum complex_abs_lt(PG_FUNCTION_ARGS)
{
    Complex *a = (Complex *) PG_GETARG_POINTER(0);
    Complex *b = (Complex *) PG_GETARG_POINTER(1);
    PG_RETURN_BOOL(complex_abs_cmp_internal(a, b) < 0);
}
```

The other four functions differ only in how they compare the internal function's result to zero.

Next we declare the functions and the operators based on the functions to SQL:

```sql
CREATE FUNCTION complex_abs_lt(complex, complex) RETURNS bool
    AS 'filename', 'complex_abs_lt'
    LANGUAGE C IMMUTABLE STRICT;

CREATE OPERATOR < (
    leftarg = complex, rightarg = complex, procedure = complex_abs_lt,
    commutator = > , negator = >= ,
    restrict = scalarltsel, join = scalarltjoinsel
);```
It is important to specify the correct commutator and negator operators, as well as suitable restriction and join selectivity functions, otherwise the optimizer will be unable to make effective use of the index.

Other things worth noting are happening here:

• There can only be one operator named, say, = and taking type complex for both operands. In this case we don't have any other operator = for complex, but if we were building a practical data type we'd probably want = to be the ordinary equality operation for complex numbers (and not the equality of the absolute values). In that case, we'd need to use some other operator name for complex_abs_eq.

• Although PostgreSQL can cope with functions having the same SQL name as long as they have different argument data types, C can only cope with one global function having a given name. So we shouldn't name the C function something simple like abs_eq. Usually it's a good practice to include the data type name in the C function name, so as not to conflict with functions for other data types.

• We could have made the SQL name of the function abs_eq, relying on PostgreSQL to distinguish it by argument data types from any other SQL function of the same name. To keep the example simple, we make the function have the same names at the C level and SQL level.

The next step is the registration of the support routine required by B-trees. The example C code that implements this is in the same file that contains the operator functions. This is how we declare the function:

```sql
CREATE FUNCTION complex_abs_cmp(complex, complex)
    RETURNS integer
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;
```

Now that we have the required operators and support routine, we can finally create the operator class:

```sql
CREATE OPERATOR CLASS complex_abs_ops
    DEFAULT FOR TYPE complex USING btree AS
    OPERATOR 1 < ,
    OPERATOR 2 <= ,
    OPERATOR 3 = ,
    OPERATOR 4 >= ,
    OPERATOR 5 > ,
    FUNCTION 1 complex_abs_cmp(complex, complex);
```

And we're done! It should now be possible to create and use B-tree indexes on complex columns.

We could have written the operator entries more verbosely, as in:

```sql
    OPERATOR 1 < (complex, complex) ,
```

but there is no need to do so when the operators take the same data type we are defining the operator class for.

The above example assumes that you want to make this new operator class the default B-tree operator class for the complex data type. If you don't, just leave out the word DEFAULT.

### 38.15.5. Operator Classes and Operator Families

So far we have implicitly assumed that an operator class deals with only one data type. While there certainly can be only one data type in a particular index column, it is often useful to index operations that compare an indexed column to a value of a different data type. Also, if there is use for a cross-
data-type operator in connection with an operator class, it is often the case that the other data type has a related operator class of its own. It is helpful to make the connections between related classes explicit, because this can aid the planner in optimizing SQL queries (particularly for B-tree operator classes, since the planner contains a great deal of knowledge about how to work with them).

To handle these needs, PostgreSQL uses the concept of an operator family. An operator family contains one or more operator classes, and can also contain indexable operators and corresponding support functions that belong to the family as a whole but not to any single class within the family. We say that such operators and functions are “loose” within the family, as opposed to being bound into a specific class. Typically each operator class contains single-data-type operators while cross-data-type operators are loose in the family.

All the operators and functions in an operator family must have compatible semantics, where the compatibility requirements are set by the index method. You might therefore wonder why bother to single out particular subsets of the family as operator classes; and indeed for many purposes the class divisions are irrelevant and the family is the only interesting grouping. The reason for defining operator classes is that they specify how much of the family is needed to support any particular index. If there is an index using an operator class, then that operator class cannot be dropped without dropping the index — but other parts of the operator family, namely other operator classes and loose operators, could be dropped. Thus, an operator class should be specified to contain the minimum set of operators and functions that are reasonably needed to work with an index on a specific data type, and then related but non-essential operators can be added as loose members of the operator family.

As an example, PostgreSQL has a built-in B-tree operator family integer_ops, which includes operator classes int8_ops, int4_ops, and int2_ops for indexes on bigint (int8), integer (int4), and smallint (int2) columns respectively. The family also contains cross-data-type comparison operators allowing any two of these types to be compared, so that an index on one of these types can be searched using a comparison value of another type. The family could be duplicated by these definitions:

CREATE OPERATOR FAMILY integer_ops USING btree;

CREATE OPERATOR CLASS int8_ops DEFAULT FOR TYPE int8 USING btree FAMILY integer_ops AS
  -- standard int8 comparisons
  OPERATOR 1 < ,
  OPERATOR 2 <= ,
  OPERATOR 3 = ,
  OPERATOR 4 >= ,
  OPERATOR 5 > ,
  FUNCTION 1 btint8cmp(int8, int8) ,
  FUNCTION 2 btint8sortsupport(internal) ,
  FUNCTION 3 in_range(int8, int8, int8, boolean, boolean) ;

CREATE OPERATOR CLASS int4_ops DEFAULT FOR TYPE int4 USING btree FAMILY integer_ops AS
  -- standard int4 comparisons
  OPERATOR 1 < ,
  OPERATOR 2 <= ,
  OPERATOR 3 = ,
  OPERATOR 4 >= ,
  OPERATOR 5 > ,
  FUNCTION 1 btint4cmp(int4, int4) ,
  FUNCTION 2 btint4sortsupport(internal) ,
  FUNCTION 3 in_range(int4, int4, int4, boolean, boolean) ;

CREATE OPERATOR CLASS int2_ops DEFAULT FOR TYPE int2 USING btree FAMILY integer_ops AS
Extending SQL

-- standard int2 comparisons
OPERATOR 1 < ,
OPERATOR 2 <= ,
OPERATOR 3 = ,
OPERATOR 4 >= ,
OPERATOR 5 > ,
FUNCTION 1 btint2cmp(int2, int2) ,
FUNCTION 2 btint2sortsupport(internal) ,
FUNCTION 3 in_range(int2, int2, int2, boolean, boolean) ;

ALTER OPERATOR FAMILY integer_ops USING btree ADD
-- cross-type comparisons int8 vs int2
OPERATOR 1 < (int8, int2) ,
OPERATOR 2 <= (int8, int2) ,
OPERATOR 3 = (int8, int2) ,
OPERATOR 4 >= (int8, int2) ,
OPERATOR 5 > (int8, int2) ,
FUNCTION 1 btint82cmp(int8, int2) ,

-- cross-type comparisons int8 vs int4
OPERATOR 1 < (int8, int4) ,
OPERATOR 2 <= (int8, int4) ,
OPERATOR 3 = (int8, int4) ,
OPERATOR 4 >= (int8, int4) ,
OPERATOR 5 > (int8, int4) ,
FUNCTION 1 btint84cmp(int8, int4) ,

-- cross-type comparisons int4 vs int2
OPERATOR 1 < (int4, int2) ,
OPERATOR 2 <= (int4, int2) ,
OPERATOR 3 = (int4, int2) ,
OPERATOR 4 >= (int4, int2) ,
OPERATOR 5 > (int4, int2) ,
FUNCTION 1 btint42cmp(int4, int2) ,

-- cross-type comparisons int4 vs int8
OPERATOR 1 < (int4, int8) ,
OPERATOR 2 <= (int4, int8) ,
OPERATOR 3 = (int4, int8) ,
OPERATOR 4 >= (int4, int8) ,
OPERATOR 5 > (int4, int8) ,
FUNCTION 1 btint48cmp(int4, int8) ,

-- cross-type comparisons int2 vs int8
OPERATOR 1 < (int2, int8) ,
OPERATOR 2 <= (int2, int8) ,
OPERATOR 3 = (int2, int8) ,
OPERATOR 4 >= (int2, int8) ,
OPERATOR 5 > (int2, int8) ,
FUNCTION 1 btint28cmp(int2, int8) ,

-- cross-type comparisons int2 vs int4
OPERATOR 1 < (int2, int4) ,
OPERATOR 2 <= (int2, int4) ,
OPERATOR 3 = (int2, int4) ,
OPERATOR 4 >= (int2, int4) ,
OPERATOR 5 > (int2, int4) ,
FUNCTION 1 btint24cmp(int2, int4) ,
-- cross-type in_range functions
FUNCTION 3 in_range(int4, int4, int8, boolean, boolean) ,
FUNCTION 3 in_range(int4, int4, int2, boolean, boolean) ,
FUNCTION 3 in_range(int2, int2, int8, boolean, boolean) ,
FUNCTION 3 in_range(int2, int2, int4, boolean, boolean) ;

Notice that this definition “overloads” the operator strategy and support function numbers: each number occurs multiple times within the family. This is allowed so long as each instance of a particular number has distinct input data types. The instances that have both input types equal to an operator class’s input type are the primary operators and support functions for that operator class, and in most cases should be declared as part of the operator class rather than as loose members of the family.

In a B-tree operator family, all the operators in the family must sort compatibly, as is specified in detail in Section 63.2. For each operator in the family there must be a support function having the same two input data types as the operator. It is recommended that a family be complete, i.e., for each combination of data types, all operators are included. Each operator class should include just the non-cross-type operators and support function for its data type.

To build a multiple-data-type hash operator family, compatible hash support functions must be created for each data type supported by the family. Here compatibility means that the functions are guaranteed to return the same hash code for any two values that are considered equal by the family’s equality operators, even when the values are of different types. This is usually difficult to accomplish when the types have different physical representations, but it can be done in some cases. Furthermore, casting a value from one data type represented in the operator family to another data type also represented in the operator family via an implicit or binary coercion cast must not change the computed hash value. Notice that there is only one support function per data type, not one per equality operator. It is recommended that a family be complete, i.e., provide an equality operator for each combination of data types. Each operator class should include just the non-cross-type equality operator and the support function for its data type.

GiST, SP-GiST, and GIN indexes do not have any explicit notion of cross-data-type operations. The set of operators supported is just whatever the primary support functions for a given operator class can handle.

In BRIN, the requirements depend on the framework that provides the operator classes. For operator classes based on minmax, the behavior required is the same as for B-tree operator families: all the operators in the family must sort compatibly, and casts must not change the associated sort ordering.

Note
Prior to PostgreSQL 8.3, there was no concept of operator families, and so any cross-data-type operators intended to be used with an index had to be bound directly into the index’s operator class. While this approach still works, it is deprecated because it makes an index’s dependencies too broad, and because the planner can handle cross-data-type comparisons more effectively when both data types have operators in the same operator family.

38.15.6. System Dependencies on Operator Classes

PostgreSQL uses operator classes to infer the properties of operators in more ways than just whether they can be used with indexes. Therefore, you might want to create operator classes even if you have no intention of indexing any columns of your data type.

In particular, there are SQL features such as ORDER BY and DISTINCT that require comparison and sorting of values. To implement these features on a user-defined data type, PostgreSQL looks for the default B-tree operator class for the data type. The “equals” member of this operator class defines the
system's notion of equality of values for `GROUP BY` and `DISTINCT`, and the sort ordering imposed by the operator class defines the default `ORDER BY` ordering.

If there is no default B-tree operator class for a data type, the system will look for a default hash operator class. But since that kind of operator class only provides equality, it is only able to support grouping not sorting.

When there is no default operator class for a data type, you will get errors like "could not identify an ordering operator" if you try to use these SQL features with the data type.

**Note**

In PostgreSQL versions before 7.4, sorting and grouping operations would implicitly use operators named `=`, `<`, and `>`. The new behavior of relying on default operator classes avoids having to make any assumption about the behavior of operators with particular names.

Sorting by a non-default B-tree operator class is possible by specifying the class's less-than operator in a `USING` option, for example

```sql
SELECT * FROM mytable ORDER BY somecol USING ~<<;
```

Alternatively, specifying the class's greater-than operator in `USING` selects a descending-order sort.

Comparison of arrays of a user-defined type also relies on the semantics defined by the type's default B-tree operator class. If there is no default B-tree operator class, but there is a default hash operator class, then array equality is supported, but not ordering comparisons.

Another SQL feature that requires even more data-type-specific knowledge is the `RANGE` offset `PRECEDING/FOLLOWING` framing option for window functions (see Section 4.2.8). For a query such as

```sql
SELECT sum(x) OVER (ORDER BY x RANGE BETWEEN 5 PRECEDING AND 10 FOLLOWING) FROM mytable;
```

it is not sufficient to know how to order by `x`; the database must also understand how to "subtract 5" or "add 10" to the current row's value of `x` to identify the bounds of the current window frame. Comparing the resulting bounds to other rows' values of `x` is possible using the comparison operators provided by the B-tree operator class that defines the `ORDER BY` ordering — but addition and subtraction operators are not part of the operator class, so which ones should be used? Hard-wiring that choice would be undesirable, because different sort orders (different B-tree operator classes) might need different behavior. Therefore, a B-tree operator class can specify an `in_range` support function that encapsulates the addition and subtraction behaviors that make sense for its sort order. It can even provide more than one `in_range` support function, in case there is more than one data type that makes sense to use as the offset in `RANGE` clauses. If the B-tree operator class associated with the window's `ORDER BY` clause does not have a matching `in_range` support function, the `RANGE offset PRECEDING/FOLLOWING` option is not supported.

Another important point is that an equality operator that appears in a hash operator family is a candidate for hash joins, hash aggregation, and related optimizations. The hash operator family is essential here since it identifies the hash function(s) to use.

### 38.15.7. Ordering Operators

Some index access methods (currently, only GiST) support the concept of ordering operators. What we have been discussing so far are search operators. A search operator is one for which the index can
be searched to find all rows satisfying \texttt{WHERE indexed\_column \ operator constant}. Note that nothing is promised about the order in which the matching rows will be returned. In contrast, an ordering operator does not restrict the set of rows that can be returned, but instead determines their order. An ordering operator is one for which the index can be scanned to return rows in the order represented by \texttt{ORDER BY indexed\_column \ operator constant}. The reason for defining ordering operators that way is that it supports nearest-neighbor searches, if the operator is one that measures distance. For example, a query like

\begin{verbatim}
SELECT * FROM places ORDER BY location <-> point '(101,456)' LIMIT 10;
\end{verbatim}

finds the ten places closest to a given target point. A GiST index on the location column can do this efficiently because <-> is an ordering operator.

While search operators have to return Boolean results, ordering operators usually return some other type, such as float or numeric for distances. This type is normally not the same as the data type being indexed. To avoid hard-wiring assumptions about the behavior of different data types, the definition of an ordering operator is required to name a B-tree operator family that specifies the sort ordering of the result data type. As was stated in the previous section, B-tree operator families define PostgreSQL's notion of ordering, so this is a natural representation. Since the point <-> operator returns float8, it could be specified in an operator class creation command like this:

\begin{verbatim}
OPERATOR 15    <-> (point, point) FOR ORDER BY float\_ops
\end{verbatim}

where float\_ops is the built-in operator family that includes operations on float8. This declaration states that the index is able to return rows in order of increasing values of the <-> operator.

### 38.15.8. Special Features of Operator Classes

There are two special features of operator classes that we have not discussed yet, mainly because they are not useful with the most commonly used index methods.

Normally, declaring an operator as a member of an operator class (or family) means that the index method can retrieve exactly the set of rows that satisfy a \texttt{WHERE} condition using the operator. For example:

\begin{verbatim}
SELECT * FROM table WHERE integer\_column < 4;
\end{verbatim}

can be satisfied exactly by a B-tree index on the integer column. But there are cases where an index is useful as an inexact guide to the matching rows. For example, if a GiST index stores only bounding boxes for geometric objects, then it cannot exactly satisfy a \texttt{WHERE} condition that tests overlap between nonrectangular objects such as polygons. Yet we could use the index to find objects whose bounding box overlaps the bounding box of the target object, and then do the exact overlap test only on the objects found by the index. If this scenario applies, the index is said to be "lossy" for the operator. Lossy index searches are implemented by having the index method return a recheck flag when a row might or might not really satisfy the query condition. The core system will then test the original query condition on the retrieved row to see whether it should be returned as a valid match. This approach works if the index is guaranteed to return all the required rows, plus perhaps some additional rows, which can be eliminated by performing the original operator invocation. The index methods that support lossy searches (currently, GiST, SP-GiST and GIN) allow the support functions of individual operator classes to set the recheck flag, and so this is essentially an operator-class feature.

Consider again the situation where we are storing in the index only the bounding box of a complex object such as a polygon. In this case there's not much value in storing the whole polygon in the index entry — we might as well store just a simpler object of type box. This situation is expressed by the \texttt{STORAGE} option in \texttt{CREATE OPERATOR CLASS}: we'd write something like:
CREATE OPERATOR CLASS polygon_ops
  DEFAULT FOR TYPE polygon USING gist AS
  ...;

At present, only the GiST, GIN and BRIN index methods support a STORAGE type that’s different from the column data type. The GiST compress and decompress support routines must deal with data-type conversion when STORAGE is used. In GIN, the STORAGE type identifies the type of the “key” values, which normally is different from the type of the indexed column — for example, an operator class for integer-array columns might have keys that are just integers. The GIN extract-Value and extractQuery support routines are responsible for extracting keys from indexed values. BRIN is similar to GIN: the STORAGE type identifies the type of the stored summary values, and operator classes’ support procedures are responsible for interpreting the summary values correctly.

38.16. Packaging Related Objects into an Extension

A useful extension to PostgreSQL typically includes multiple SQL objects; for example, a new data type will require new functions, new operators, and probably new index operator classes. It is helpful to collect all these objects into a single package to simplify database management. PostgreSQL calls such a package an extension. To define an extension, you need at least a script file that contains the SQL commands to create the extension’s objects, and a control file that specifies a few basic properties of the extension itself. If the extension includes C code, there will typically also be a shared library file into which the C code has been built. Once you have these files, a simple CREATE EXTENSION command loads the objects into your database.

The main advantage of using an extension, rather than just running the SQL script to load a bunch of “loose” objects into your database, is that PostgreSQL will then understand that the objects of the extension go together. You can drop all the objects with a single DROP EXTENSION command (no need to maintain a separate “uninstall” script). Even more useful, pg_dump knows that it should not dump the individual member objects of the extension — it will just include a CREATE EXTENSION command in dumps, instead. This vastly simplifies migration to a new version of the extension that might contain more or different objects than the old version. Note however that you must have the extension’s control, script, and other files available when loading such a dump into a new database.

PostgreSQL will not let you drop an individual object contained in an extension, except by dropping the whole extension. Also, while you can change the definition of an extension member object (for example, via CREATE OR REPLACE FUNCTION for a function), bear in mind that the modified definition will not be dumped by pg_dump. Such a change is usually only sensible if you concurrently make the same change in the extension’s script file. (But there are special provisions for tables containing configuration data; see Section 38.16.4.) In production situations, it’s generally better to create an extension update script to perform changes to extension member objects.

The extension script may set privileges on objects that are part of the extension via GRANT and REVOKE statements. The final set of privileges for each object (if any are set) will be stored in the pg_init_privs system catalog. When pg_dump is used, the CREATE EXTENSION command will be included in the dump, followed by the set of GRANT and REVOKE statements necessary to set the privileges on the objects to what they were at the time the dump was taken.

PostgreSQL does not currently support extension scripts issuing CREATE POLICY or SECURITY LABEL statements. These are expected to be set after the extension has been created. All RLS policies and security labels on extension objects will be included in dumps created by pg_dump.

The extension mechanism also has provisions for packaging modification scripts that adjust the definitions of the SQL objects contained in an extension. For example, if version 1.1 of an extension adds one function and changes the body of another function compared to 1.0, the extension author can provide an update script that makes just those two changes. The ALTER EXTENSION UPDATE
command can then be used to apply these changes and track which version of the extension is actually
installed in a given database.

The kinds of SQL objects that can be members of an extension are shown in the description of AL-
TER EXTENSION. Notably, objects that are database-cluster-wide, such as databases, roles, and ta-
blespaces, cannot be extension members since an extension is only known within one database. (Al-
though an extension script is not prohibited from creating such objects, if it does so they will not be
tracked as part of the extension.) Also notice that while a table can be a member of an extension, its
subsidiary objects such as indexes are not directly considered members of the extension. Another im-
portant point is that schemas can belong to extensions, but not vice versa: an extension as such has an
unqualified name and does not exist “within” any schema. The extension’s member objects, however,
will belong to schemas whenever appropriate for their object types. It may or may not be appropriate
for an extension to own the schema(s) its member objects are within.

If an extension’s script creates any temporary objects (such as temp tables), those objects are treated as
extension members for the remainder of the current session, but are automatically dropped at session
end, as any temporary object would be. This is an exception to the rule that extension member objects
cannot be dropped without dropping the whole extension.

38.16.1. Defining Extension Objects

Widely-distributed extensions should assume little about the database they occupy. In particular, un-
less you issued SET search_path = pg_temp, assume each unqualified name could resolve
to an object that a malicious user has defined. Beware of constructs that depend on search_path
implicitly: IN and CASE expression WHEN always select an operator using the search path. In
their place, use OPERATOR(schema.=) ANY and CASE WHEN expression.

38.16.2. Extension Files

The CREATE EXTENSION command relies on a control file for each extension, which must be
named the same as the extension with a suffix of .control, and must be placed in the installation's
SHAREDIR/extension directory. There must also be at least one SQL script file, which follows
the naming pattern extension--version.sql (for example, foo--1.0.sql for version 1.0
of extension foo). By default, the script file(s) are also placed in the SHAREDIR/extension di-
rectory; but the control file can specify a different directory for the script file(s).

The file format for an extension control file is the same as for the postgresql.conf file, namely a
list of parameter_name = value assignments, one per line. Blank lines and comments introduced
by # are allowed. Be sure to quote any value that is not a single word or number.

A control file can set the following parameters:

directory(string)

The directory containing the extension’s SQL script file(s). Unless an absolute path is given, the
name is relative to the installation's SHAREDIR directory. The default behavior is equivalent to
specifying directory = 'extension'.

default_version(string)

The default version of the extension (the one that will be installed if no version is specified in
CREATE EXTENSION). Although this can be omitted, that will result in CREATE EXTENSION
failing if no VERSION option appears, so you generally don't want to do that.

comment(string)

A comment (any string) about the extension. The comment is applied when initially creating an
extension, but not during extension updates (since that might override user-added comments).
Alternatively, the extension’s comment can be set by writing a COMMENT command in the script
file.
encoding(string)

The character set encoding used by the script file(s). This should be specified if the script files contain any non-ASCII characters. Otherwise the files will be assumed to be in the database encoding.

module_pathname(string)

The value of this parameter will be substituted for each occurrence of MODULE_PATHNAME in the script file(s). If it is not set, no substitution is made. Typically, this is set to $lib-dir/shared_library_name and then MODULE_PATHNAME is used in CREATE FUNCTION commands for C-language functions, so that the script files do not need to hard-wire the name of the shared library.

requires(string)

A list of names of extensions that this extension depends on, for example requires = 'foo, bar'. Those extensions must be installed before this one can be installed.

superuser(boolean)

If this parameter is true (which is the default), only superusers can create the extension or update it to a new version. If it is set to false, just the privileges required to execute the commands in the installation or update script are required.

relocatable(boolean)

An extension is relocatable if it is possible to move its contained objects into a different schema after initial creation of the extension. The default is false, i.e. the extension is not relocatable. See Section 38.16.3 for more information.

schema(string)

This parameter can only be set for non-relocatable extensions. It forces the extension to be loaded into exactly the named schema and not any other. The schema parameter is consulted only when initially creating an extension, not during extension updates. See Section 38.16.3 for more information.

In addition to the primary control file extension.control, an extension can have secondary control files named in the style extension--version.control. If supplied, these must be located in the script file directory. Secondary control files follow the same format as the primary control file. Any parameters set in a secondary control file override the primary control file when installing or updating to that version of the extension. However, the parameters directory and default_version cannot be set in a secondary control file.

An extension's SQL script files can contain any SQL commands, except for transaction control commands (BEGIN, COMMIT, etc) and commands that cannot be executed inside a transaction block (such as VACUUM). This is because the script files are implicitly executed within a transaction block.

An extension's SQL script files can also contain lines beginning with \echo, which will be ignored (treated as comments) by the extension mechanism. This provision is commonly used to throw an error if the script file is fed to psql rather than being loaded via CREATE EXTENSION (see example script in Section 38.16.7). Without that, users might accidentally load the extension's contents as “loose” objects rather than as an extension, a state of affairs that's a bit tedious to recover from.

While the script files can contain any characters allowed by the specified encoding, control files should contain only plain ASCII, because there is no way for PostgreSQL to know what encoding a control file is in. In practice this is only an issue if you want to use non-ASCII characters in the extension's comment. Recommended practice in that case is to not use the control file comment parameter, but instead use COMMENT ON EXTENSION within a script file to set the comment.
38.16.3. Extension Relocatability

Users often wish to load the objects contained in an extension into a different schema than the extension's author had in mind. There are three supported levels of relocatability:

- A fully relocatable extension can be moved into another schema at any time, even after it's been loaded into a database. This is done with the `ALTER EXTENSION SET SCHEMA` command, which automatically renames all the member objects into the new schema. Normally, this is only possible if the extension contains no internal assumptions about what schema any of its objects are in. Also, the extension's objects must all be in one schema to begin with (ignoring objects that do not belong to any schema, such as procedural languages). Mark a fully relocatable extension by setting `relocatable = true` in its control file.

- An extension might be relocatable during installation but not afterwards. This is typically the case if the extension's script file needs to reference the target schema explicitly, for example in setting `search_path` properties for SQL functions. For such an extension, set `relocatable = false` in its control file, and use `@extschema@` to refer to the target schema in the script file. All occurrences of this string will be replaced by the actual target schema's name before the script is executed. The user can set the target schema using the `SCHEMA` option of `CREATE EXTENSION`.

- If the extension does not support relocation at all, set `relocatable = false` in its control file, and also set `schema` to the name of the intended target schema. This will prevent use of the `SCHEMA` option of `CREATE EXTENSION`, unless it specifies the same schema named in the control file. This choice is typically necessary if the extension contains internal assumptions about schema names that can't be replaced by uses of `@extschema@`. The `@extschema@` substitution mechanism is available in this case too, although it is of limited use since the schema name is determined by the control file.

In all cases, the script file will be executed with `search_path` initially set to point to the target schema; that is, `CREATE EXTENSION` does the equivalent of this:

```
SET LOCAL search_path TO @extschema@;
```

This allows the objects created by the script file to go into the target schema. The script file can change `search_path` if it wishes, but that is generally undesirable. `search_path` is restored to its previous setting upon completion of `CREATE EXTENSION`.

The target schema is determined by the `schema` parameter in the control file if that is given, otherwise by the `SCHEMA` option of `CREATE EXTENSION` if that is given, otherwise the current default object creation schema (the first one in the caller's `search_path`). When the control file `schema` parameter is used, the target schema will be created if it doesn't already exist, but in the other two cases it must already exist.

If any prerequisite extensions are listed in `requires` in the control file, their target schemas are appended to the initial setting of `search_path`. This allows their objects to be visible to the new extension's script file.

Although a non-relocatable extension can contain objects spread across multiple schemas, it is usually desirable to place all the objects meant for external use into a single schema, which is considered the extension's target schema. Such an arrangement works conveniently with the default setting of `search_path` during creation of dependent extensions.

38.16.4. Extension Configuration Tables

Some extensions include configuration tables, which contain data that might be added or changed by the user after installation of the extension. Ordinarily, if a table is part of an extension, neither the table's definition nor its content will be dumped by `pg_dump`. But that behavior is undesirable for a configuration table; any data changes made by the user need to be included in dumps, or the extension will behave differently after a dump and reload.
To solve this problem, an extension's script file can mark a table or a sequence it has created as a configuration relation, which will cause pg_dump to include the table's or the sequence's contents (not its definition) in dumps. To do that, call the function pg_extension_config_dump(regclass, text) after creating the table or the sequence, for example

```
CREATE TABLE my_config (key text, value text);
CREATE SEQUENCE my_config_seq;
```

```
SELECT pg_catalog.pg_extension_config_dump('my_config', '');
SELECT pg_catalog.pg_extension_config_dump('my_config_seq', '');
```

Any number of tables or sequences can be marked this way. Sequences associated with `serial` or `bigserial` columns can be marked as well.

When the second argument of pg_extension_config_dump is an empty string, the entire contents of the table are dumped by pg_dump. This is usually only correct if the table is initially empty as created by the extension script. If there is a mixture of initial data and user-provided data in the table, the second argument of pg_extension_config_dump provides a WHERE condition that selects the data to be dumped. For example, you might do

```
CREATE TABLE my_config (key text, value text, standard_entry boolean);
SELECT pg_catalog.pg_extension_config_dump('my_config', 'WHERE NOT standard_entry');
```

and then make sure that `standard_entry` is true only in the rows created by the extension's script.

For sequences, the second argument of pg_extension_config_dump has no effect.

More complicated situations, such as initially-provided rows that might be modified by users, can be handled by creating triggers on the configuration table to ensure that modified rows are marked correctly.

You can alter the filter condition associated with a configuration table by calling pg_extension_config_dump again. (This would typically be useful in an extension update script.) The only way to mark a table as no longer a configuration table is to dissociate it from the extension with ALTER EXTENSION ... DROP TABLE.

Note that foreign key relationships between these tables will dictate the order in which the tables are dumped out by pg_dump. Specifically, pg_dump will attempt to dump the referenced-by table before the referencing table. As the foreign key relationships are set up at CREATE EXTENSION time (prior to data being loaded into the tables) circular dependencies are not supported. When circular dependencies exist, the data will still be dumped out but the dump will not be able to be restored directly and user intervention will be required.

Sequences associated with `serial` or `bigserial` columns need to be directly marked to dump their state. Marking their parent relation is not enough for this purpose.

### 38.16.5. Extension Updates

One advantage of the extension mechanism is that it provides convenient ways to manage updates to the SQL commands that define an extension's objects. This is done by associating a version name or number with each released version of the extension's installation script. In addition, if you want users to be able to update their databases dynamically from one version to the next, you should provide update scripts that make the necessary changes to go from one version to the next. Update scripts have names following the pattern `extension--oldversion--newversion.sql` (for exam-
Extending SQL

Example, foo--1.0--1.1.sql contains the commands to modify version 1.0 of extension foo into version 1.1).

Given that a suitable update script is available, the command ALTER EXTENSION UPDATE will update an installed extension to the specified new version. The update script is run in the same environment that CREATE EXTENSION provides for installation scripts: in particular, search_path is set up in the same way, and any new objects created by the script are automatically added to the extension. Also, if the script chooses to drop extension member objects, they are automatically disassociated from the extension.

If an extension has secondary control files, the control parameters that are used for an update script are those associated with the script’s target (new) version.

The update mechanism can be used to solve an important special case: converting a “loose” collection of objects into an extension. Before the extension mechanism was added to PostgreSQL (in 9.1), many people wrote extension modules that simply created assorted unpackaged objects. Given an existing database containing such objects, how can we convert the objects into a properly packaged extension? Dropping them and then doing a plain CREATE EXTENSION is one way, but it’s not desirable if the objects have dependencies (for example, if there are table columns of a data type created by the extension). The way to fix this situation is to create an empty extension, then use ALTER EXTENSION ADD to attach each pre-existing object to the extension, then finally create any new objects that are in the current extension version but were not in the unpackaged release. CREATE EXTENSION supports this case with its FROM old_version option, which causes it to not run the normal installation script for the target version, but instead the update script named extension--old_version--target_version.sql. The choice of the dummy version name to use as old_version is up to the extension author, though unpackaged is a common convention. If you have multiple prior versions you need to be able to update into extension style, use multiple dummy version names to identify them.

ALTER EXTENSION is able to execute sequences of update script files to achieve a requested update. For example, if only foo--1.0--1.1.sql and foo--1.1--2.0.sql are available, ALTER EXTENSION will apply them in sequence if an update to version 2.0 is requested when 1.0 is currently installed.

PostgreSQL doesn’t assume anything about the properties of version names: for example, it does not know whether 1.1 follows 1.0. It just matches up the available version names and follows the path that requires applying the fewest update scripts. (A version name can actually be any string that doesn’t contain -- or leading or trailing --.)

Sometimes it is useful to provide “downgrade” scripts, for example foo--1.1--1.0.sql to allow reverting the changes associated with version 1.1. If you do that, be careful of the possibility that a downgrade script might unexpectedly get applied because it yields a shorter path. The risky case is where there is a “fast path” update script that jumps ahead several versions as well as a downgrade script to the fast path’s start point. It might take fewer steps to apply the downgrade and then the fast path than to move ahead one version at a time. If the downgrade script drops any irreplaceable objects, this will yield undesirable results.

To check for unexpected update paths, use this command:

```sql
SELECT * FROM pg_extension_update_paths('extension_name');
```

This shows each pair of distinct known version names for the specified extension, together with the update path sequence that would be taken to get from the source version to the target version, or NULL if there is no available update path. The path is shown in textual form with -- separators. You can use regexp_split_to_array(path, '--') if you prefer an array format.

### 38.16.6. Installing Extensions using Update Scripts

An extension that has been around for awhile will probably exist in several versions, for which the author will need to write update scripts. For example, if you have released a foo exten-
sion in versions 1.0, 1.1, and 1.2, there should be update scripts foo--1.0--1.1.sql and foo--1.1--1.2.sql. Before PostgreSQL 10, it was necessary to also create new script files foo--1.1.sql and foo--1.2.sql that directly build the newer extension versions, or else the newer versions could not be installed directly, only by installing 1.0 and then updating. That was tedious and duplicative, but now it's unnecessary, because CREATE EXTENSION can follow update chains automatically. For example, if only the script files foo--1.0.sql, foo--1.0--1.1.sql, and foo--1.1--1.2.sql are available then a request to install version 1.2 is honored by running those three scripts in sequence. The processing is the same as if you’d first installed 1.0 and then updated to 1.2. (As with ALTER EXTENSION UPDATE, if multiple pathways are available then the shortest is preferred.) Arranging an extension's script files in this style can reduce the amount of maintenance effort needed to produce small updates.

If you use secondary (version-specific) control files with an extension maintained in this style, keep in mind that each version needs a control file even if it has no stand-alone installation script, as that control file will determine how the implicit update to that version is performed. For example, if foo--1.0.control specifies requires = 'bar' but foo's other control files do not, the extension's dependency on bar will be dropped when updating from 1.0 to another version.

38.16.7. Extension Example

Here is a complete example of an SQL-only extension, a two-element composite type that can store any type of value in its slots, which are named “k” and “v”. Non-text values are automatically coerced to text for storage.

The script file pair--1.0.sql looks like this:

```sql
-- complain if script is sourced in psql, rather than via CREATE EXTENSION
\echo Use "CREATE EXTENSION pair" to load this file. \quit
CREATE TYPE pair AS ( k text, v text );
CREATE OR REPLACE FUNCTION pair(text, text)
RETURNS pair LANGUAGE SQL AS 'SELECT ROW($1, $2)::@extschema@.pair;';
CREATE OPERATOR ~> (LEFTARG = text, RIGHTARG = text, FUNCTION = pair);
-- "SET search_path" is easy to get right, but qualified names perform better.
CREATE OR REPLACE FUNCTION lower(pair)
RETURNS pair LANGUAGE SQL
AS 'SELECT ROW(lower($1.k), lower($1.v))::@extschema@.pair;'
SET search_path = pg_temp;
CREATE OR REPLACE FUNCTION pair_concat(pair, pair)
RETURNS pair LANGUAGE SQL
AS 'SELECT ROW($1.k OPERATOR(pg_catalog.||) $2.k, $1.v OPERATOR(pg_catalog.||) $2.v)::@extschema@.pair;';
```

The control file pair.control looks like this:

```control
# pair extension
```
While you hardly need a makefile to install these two files into the correct directory, you could use a makefile containing this:

```makefile
EXTENSION = pair
DATA = pair--1.0.sql
PG_CONFIG = pg_config
PGXS := $(shell $(PG_CONFIG) --pgxs)
include $(PGXS)
```

This makefile relies on PGXS, which is described in Section 38.17. The command `make install` will install the control and script files into the correct directory as reported by `pg_config`.

Once the files are installed, use the `CREATE EXTENSION` command to load the objects into any particular database.

### 38.17. Extension Building Infrastructure

If you are thinking about distributing your PostgreSQL extension modules, setting up a portable build system for them can be fairly difficult. Therefore the PostgreSQL installation provides a build infrastructure for extensions, called PGXS, so that simple extension modules can be built simply against an already installed server. PGXS is mainly intended for extensions that include C code, although it can be used for pure-SQL extensions too. Note that PGXS is not intended to be a universal build system framework that can be used to build any software interfacing to PostgreSQL; it simply automates common build rules for simple server extension modules. For more complicated packages, you might need to write your own build system.

To use the PGXS infrastructure for your extension, you must write a simple makefile. In the makefile, you need to set some variables and include the global PGXS makefile. Here is an example that builds an extension module named `isbn_issn`, consisting of a shared library containing some C code, an extension control file, a SQL script, an include file (only needed if other modules might need to access the extension functions without going via SQL), and a documentation text file:

```makefile
MODULES = isbn_issn
EXTENSION = isbn_issn
DATA = isbn_issn--1.0.sql
DOCS = README.isbn_issn
HEADERS_isbn_issn = isbn_issn.h
PG_CONFIG = pg_config
PGXS := $(shell $(PG_CONFIG) --pgxs)
include $(PGXS)
```

The last three lines should always be the same. Earlier in the file, you assign variables or add custom make rules.

Set one of these three variables to specify what is built:

```makefile
MODULES
```

- list of shared-library objects to be built from source files with same stem (do not include library suffixes in this list)
Extending SQL

MODULE_big
a shared library to build from multiple source files (list object files in OBJS)

PROGRAM
an executable program to build (list object files in OBJS)

The following variables can also be set:

EXTENSION
extension name(s); for each name you must provide an extension.control file, which will
be installed into prefix/share/extension

MODULEDIR
subdirectory of prefix/share into which DATA and DOCS files should be installed (if not
set, default is extension if EXTENSION is set, or contrib if not)

DATA
random files to install into prefix/share/$MODULEDIR

DATA_built
random files to install into prefix/share/$MODULEDIR, which need to be built first

DATA_TSEARCH
random files to install under prefix/share/tsearch_data

DOCS
random files to install under prefix/doc/$MODULEDIR

HEADERS
HEADERS_built
Files to (optionally build and) install under prefix/include/server/$MODULEDIR/$MODULE_big.

Unlike DATA_built, files in HEADERS_built are not removed by the clean target; if you
want them removed, also add them to EXTRA_CLEAN or add your own rules to do it.

HEADERS_$MODULE
HEADERS_built_$MODULE
Files to install (after building if specified) under prefix/include/server/$MODULEDIR/$MODULE, where $MODULE must be a module name used in MODULES or MOD-
ULE_big.

Unlike DATA_built, files in HEADERS_built_$MODULE are not removed by the clean
target; if you want them removed, also add them to EXTRA_CLEAN or add your own rules to do it.

It is legal to use both variables for the same module, or any combination, unless you have
two module names in the MODULES list that differ only by the presence of a prefix built_.
which would cause ambiguity. In that (hopefully unlikely) case, you should use only the HEAD-
ERS_built_$MODULE variables.

SCRIPTS
script files (not binaries) to install into prefix/bin
Extending SQL

SCRIPTS_built
script files (not binaries) to install into prefix/bin, which need to be built first

REGRESS
list of regression test cases (without suffix), see below

REGRESS_OPTS
additional switches to pass to pg_regress

NO_INSTALLCHECK
don’t define an installcheck target, useful e.g. if tests require special configuration, or don’t use pg_regress

EXTRA_CLEAN
extra files to remove in make clean

PG_CPPFLAGS
will be added to CPPFLAGS

PG_LIBS
will be added to PROGRAM link line

SHLIB_LINK
will be added to MODULE_big link line

PG_CONFIG
path to pg_config program for the PostgreSQL installation to build against (typically just pg_config to use the first one in your PATH)

Put this makefile as Makefile in the directory which holds your extension. Then you can do make to compile, and then make install to install your module. By default, the extension is compiled and installed for the PostgreSQL installation that corresponds to the first pg_config program found in your PATH. You can use a different installation by setting PG_CONFIG to point to its pg_config program, either within the makefile or on the make command line.

You can also run make in a directory outside the source tree of your extension, if you want to keep the build directory separate. This procedure is also called a VPATH build. Here’s how:

mkdir build_dir
cd build_dir
make -f /path/to/extension/source/tree/Makefile
make -f /path/to/extension/source/tree/Makefile install

Alternatively, you can set up a directory for a VPATH build in a similar way to how it is done for the core code. One way to do this is using the core script config/prep_buildtree. Once this has been done you can build by setting the make variable VPATH like this:

make VPATH=/path/to/extension/source/tree
make VPATH=/path/to/extension/source/tree install

This procedure can work with a greater variety of directory layouts.
The scripts listed in the REGRESS variable are used for regression testing of your module, which can be invoked by make installcheck after doing make install. For this to work you must have a running PostgreSQL server. The script files listed in REGRESS must appear in a subdirectory named sql/ in your extension’s directory. These files must have extension .sql, which must not be included in the REGRESS list in the makefile. For each test there should also be a file containing the expected output in a subdirectory named expected/, with the same stem and extension .out. make installcheck executes each test script with psql, and compares the resulting output to the matching expected file. Any differences will be written to the file regression.diffs in diff -c format. Note that trying to run a test that is missing its expected file will be reported as “trouble”, so make sure you have all expected files.

**Tip**

The easiest way to create the expected files is to create empty files, then do a test run (which will of course report differences). Inspect the actual result files found in the results/ directory, then copy them to expected/ if they match what you expect from the test.
Chapter 39. Triggers

This chapter provides general information about writing trigger functions. Trigger functions can be written in most of the available procedural languages, including PL/pgSQL (Chapter 43), PL/Tcl (Chapter 44), PL/Perl (Chapter 45), and PL/Python (Chapter 46). After reading this chapter, you should consult the chapter for your favorite procedural language to find out the language-specific details of writing a trigger in it.

It is also possible to write a trigger function in C, although most people find it easier to use one of the procedural languages. It is not currently possible to write a trigger function in the plain SQL function language.

39.1. Overview of Trigger Behavior

A trigger is a specification that the database should automatically execute a particular function whenever a certain type of operation is performed. Triggers can be attached to tables (partitioned or not), views, and foreign tables.

On tables and foreign tables, triggers can be defined to execute either before or after any INSERT, UPDATE, or DELETE operation, either once per modified row, or once per SQL statement. UPDATE triggers can moreover be set to fire only if certain columns are mentioned in the SET clause of the UPDATE statement. Triggers can also fire for TRUNCATE statements. If a trigger event occurs, the trigger's function is called at the appropriate time to handle the event.

On views, triggers can be defined to execute instead of INSERT, UPDATE, or DELETE operations. Such INSTEAD OF triggers are fired once for each row that needs to be modified in the view. It is the responsibility of the trigger's function to perform the necessary modifications to the view's underlying base table(s) and, where appropriate, return the modified row as it will appear in the view. Triggers on views can also be defined to execute once per SQL statement, before or after INSERT, UPDATE, or DELETE operations. However, such triggers are fired only if there is also an INSTEAD OF trigger on the view. Otherwise, any statement targeting the view must be rewritten into a statement affecting its underlying base table(s), and then the triggers that will be fired are the ones attached to the base table(s).

The trigger function must be defined before the trigger itself can be created. The trigger function must be declared as a function taking no arguments and returning type trigger. (The trigger function receives its input through a specially-passed TriggerData structure, not in the form of ordinary function arguments.)

Once a suitable trigger function has been created, the trigger is established with CREATE TRIGGER. The same trigger function can be used for multiple triggers.

PostgreSQL offers both per-row triggers and per-statement triggers. With a per-row trigger, the trigger function is invoked once for each row that is affected by the statement that fired the trigger. In contrast, a per-statement trigger is invoked only once when an appropriate statement is executed, regardless of the number of rows affected by that statement. In particular, a statement that affects zero rows will still result in the execution of any applicable per-statement triggers. These two types of triggers are sometimes called row-level triggers and statement-level triggers, respectively. Triggers on TRUNCATE may only be defined at statement level, not per-row.

Triggers are also classified according to whether they fire before, after, or instead of the operation. These are referred to as BEFORE triggers, AFTER triggers, and INSTEAD OF triggers respectively. Statement-level BEFORE triggers naturally fire before the statement starts to do anything, while statement-level AFTER triggers fire at the very end of the statement. These types of triggers may be defined on tables, views, or foreign tables. Row-level BEFORE triggers fire immediately before a particular row is operated on, while row-level AFTER triggers fire at the end of the statement (but before
A statement that targets a parent table in an inheritance or partitioning hierarchy does not cause the statement-level triggers of affected child tables to be fired; only the parent table's statement-level triggers are fired. However, row-level triggers of any affected child tables will be fired.

If an `INSERT` contains an `ON CONFLICT DO UPDATE` clause, it is possible that the effects of row-level `BEFORE INSERT` triggers and row-level `BEFORE UPDATE` triggers can both be applied in a way that is apparent from the final state of the updated row, if an `EXCLUDED` column is referenced. There need not be an `EXCLUDED` column reference for both sets of row-level `BEFORE` triggers to execute, though. The possibility of surprising outcomes should be considered when there are both `BEFORE INSERT` and `BEFORE UPDATE` row-level triggers that change a row being inserted/updated (this can be problematic even if the modifications are more or less equivalent, if they're not also idempotent). Note that statement-level `UPDATE` triggers are executed when `ON CONFLICT DO UPDATE` is specified, regardless of whether or not any rows were affected by the `UPDATE` (and regardless of whether the alternative `UPDATE` path was ever taken). An `INSERT` with an `ON CONFLICT DO UPDATE` clause will execute statement-level `BEFORE INSERT` triggers first, then statement-level `BEFORE UPDATE` triggers, followed by statement-level `AFTER UPDATE` triggers and finally statement-level `AFTER INSERT` triggers.

If an `UPDATE` on a partitioned table causes a row to move to another partition, it will be performed as a `DELETE` from the original partition followed by an `INSERT` into the new partition. In this case, all row-level `BEFORE UPDATE` triggers and all row-level `BEFORE DELETE` triggers are fired on the original partition. Then all row-level `BEFORE INSERT` triggers are fired on the destination partition. The possibility of surprising outcomes should be considered when all these triggers affect the row being moved. As far as `AFTER ROW` triggers are concerned, `AFTER DELETE` and `AFTER INSERT` triggers are applied; but `AFTER UPDATE` triggers are not applied because the `UPDATE` has been converted to a `DELETE` and an `INSERT`. As far as statement-level triggers are concerned, none of the `DELETE` or `INSERT` triggers are fired, even if row movement occurs; only the `UPDATE` triggers defined on the target table used in the `UPDATE` statement will be fired.

Trigger functions invoked by per-statement triggers should always return `NULL`. Trigger functions invoked by per-row triggers can return a table row (a value of type `HeapTuple`) to the calling executor, if they choose. A row-level trigger fired before an operation has the following choices:

- It can return `NULL` to skip the operation for the current row. This instructs the executor to not perform the row-level operation that invoked the trigger (the insertion, modification, or deletion of a particular table row).
- For row-level `INSERT` and `UPDATE` triggers only, the returned row becomes the row that will be inserted or will replace the row being updated. This allows the trigger function to modify the row being inserted or updated.

A row-level `BEFORE` trigger that does not intend to cause either of these behaviors must be careful to return as its result the same row that was passed in (that is, the `NEW` row for `INSERT` and `UPDATE` triggers, the `OLD` row for `DELETE` triggers).

A row-level `INSTEAD OF` trigger should either return `NULL` to indicate that it did not modify any data from the view's underlying base tables, or it should return the view row that was passed in (the `NEW` row for `INSERT` and `UPDATE` operations, or the `OLD` row for `DELETE` operations). A nonnull return value is used to signal that the trigger performed the necessary data modifications in the view. This will cause the count of the number of rows affected by the command to be incremented. For `INSERT` and `UPDATE` operations, the trigger may modify the `NEW` row before returning it. This will change the data returned by `INSERT RETURNING` or `UPDATE RETURNING`, and is useful when the view will not show exactly the same data that was provided.

The return value is ignored for row-level triggers fired after an operation, and so they can return `NULL`.
If more than one trigger is defined for the same event on the same relation, the triggers will be fired in alphabetical order by trigger name. In the case of BEFORE and INSTEAD OF triggers, the possibly-modified row returned by each trigger becomes the input to the next trigger. If any BEFORE or INSTEAD OF trigger returns NULL, the operation is abandoned for that row and subsequent triggers are not fired (for that row).

A trigger definition can also specify a Boolean WHEN condition, which will be tested to see whether the trigger should be fired. In row-level triggers the WHEN condition can examine the old and/or new values of columns of the row. (Statement-level triggers can also have WHEN conditions, although the feature is not so useful for them.) In a BEFORE trigger, the WHEN condition is evaluated just before the function is or would be executed, so using WHEN is not materially different from testing the same condition at the beginning of the trigger function. However, in an AFTER trigger, the WHEN condition is evaluated just after the row update occurs, and it determines whether an event is queued to fire the trigger at the end of statement. So when an AFTER trigger's WHEN condition does not return true, it is not necessary to queue an event nor to re-fetch the row at end of statement. This can result in significant speedups in statements that modify many rows, if the trigger only needs to be fired for a few of the rows. INSTEAD OF triggers do not support WHEN conditions.

Typically, row-level BEFORE triggers are used for checking or modifying the data that will be inserted or updated. For example, a BEFORE trigger might be used to insert the current time into a timestamp column, or to check that two elements of the row are consistent. Row-level AFTER triggers are most sensibly used to propagate the updates to other tables, or make consistency checks against other tables. The reason for this division of labor is that an AFTER trigger can be certain it is seeing the final value of the row, while a BEFORE trigger cannot; there might be other BEFORE triggers firing after it. If you have no specific reason to make a trigger BEFORE or AFTER, the BEFORE case is more efficient, since the information about the operation doesn't have to be saved until end of statement.

If a trigger function executes SQL commands then these commands might fire triggers again. This is known as cascading triggers. There is no direct limitation on the number of cascade levels. It is possible for cascades to cause a recursive invocation of the same trigger; for example, an INSERT trigger might execute a command that inserts an additional row into the same table, causing the INSERT trigger to be fired again. It is the trigger programmer's responsibility to avoid infinite recursion in such scenarios.

When a trigger is being defined, arguments can be specified for it. The purpose of including arguments in the trigger definition is to allow different triggers with similar requirements to call the same function. As an example, there could be a generalized trigger function that takes as its arguments two column names and puts the current user in one and the current time stamp in the other. Properly written, this trigger function would be independent of the specific table it is triggering on. So the same function could be used for INSERT events on any table with suitable columns, to automatically track creation of records in a transaction table for example. It could also be used to track last-update events if defined as an UPDATE trigger.

Each programming language that supports triggers has its own method for making the trigger input data available to the trigger function. This input data includes the type of trigger event (e.g., INSERT or UPDATE) as well as any arguments that were listed in CREATE TRIGGER. For a row-level trigger, the input data also includes the NEW row for INSERT and UPDATE triggers, and/or the OLD row for UPDATE and DELETE triggers.

By default, statement-level triggers do not have any way to examine the individual row(s) modified by the statement. But an AFTER STATEMENT trigger can request that transition tables be created to make the sets of affected rows available to the trigger. AFTER ROW triggers can also request transition tables, so that they can see the total changes in the table as well as the change in the individual row they are currently being fired for. The method for examining the transition tables again depends on the programming language that is being used, but the typical approach is to make the transition tables act like read-only temporary tables that can be accessed by SQL commands issued within the trigger function.
39.2. Visibility of Data Changes

If you execute SQL commands in your trigger function, and these commands access the table that the trigger is for, then you need to be aware of the data visibility rules, because they determine whether these SQL commands will see the data change that the trigger is fired for. Briefly:

- Statement-level triggers follow simple visibility rules: none of the changes made by a statement are visible to statement-level BEFORE triggers, whereas all modifications are visible to statement-level AFTER triggers.

- The data change (insertion, update, or deletion) causing the trigger to fire is naturally not visible to SQL commands executed in a row-level BEFORE trigger, because it hasn't happened yet.

- However, SQL commands executed in a row-level BEFORE trigger will see the effects of data changes for rows previously processed in the same outer command. This requires caution, since the ordering of these change events is not in general predictable; a SQL command that affects multiple rows can visit the rows in any order.

- Similarly, a row-level INSTEAD OF trigger will see the effects of data changes made by previous firings of INSTEAD OF triggers in the same outer command.

- When a row-level AFTER trigger is fired, all data changes made by the outer command are already complete, and are visible to the invoked trigger function.

If your trigger function is written in any of the standard procedural languages, then the above statements apply only if the function is declared VOLATILE. Functions that are declared STABLE or IMMUTABLE will not see changes made by the calling command in any case.

Further information about data visibility rules can be found in Section 47.5. The example in Section 39.4 contains a demonstration of these rules.

39.3. Writing Trigger Functions in C

This section describes the low-level details of the interface to a trigger function. This information is only needed when writing trigger functions in C. If you are using a higher-level language then these details are handled for you. In most cases you should consider using a procedural language before writing your triggers in C. The documentation of each procedural language explains how to write a trigger in that language.

Trigger functions must use the “version 1” function manager interface.

When a function is called by the trigger manager, it is not passed any normal arguments, but it is passed a “context” pointer pointing to a TriggerData structure. C functions can check whether they were called from the trigger manager or not by executing the macro:

\[
\text{CALLED\_AS\_TRIGGER}(\text{fcinfo})
\]

which expands to:

\[
((\text{fcinfo})\rightarrow\text{context} \neq \text{NULL} \&\& \text{IsA((fcinfo)\rightarrow context, TriggerData)})
\]

If this returns true, then it is safe to cast \(\text{fcinfo}\rightarrow\text{context}\) to type TriggerData * and make use of the pointed-to TriggerData structure. The function must not alter the TriggerData structure or any of the data it points to.

**struct TriggerData is defined in commands/trigger.h:**
typedef struct TriggerData {
    NodeTag type;
    TriggerEvent tg_event;
    Relation tg_relation;
    HeapTuple tg_trigtuple;
    HeapTuple tg_newtuple;
    Trigger *tg_trigger;
    Buffer tg_trigtuplebuf;
    Buffer tg_newtuplebuf;
    Tuplestorestate *tg_oldtable;
    Tuplestorestate *tg_newtable;
} TriggerData;

where the members are defined as follows:

type

Always T_TriggerData.

tg_event

Describes the event for which the function is called. You can use the following macros to examine tg_event:

TRIGGER_FIRED_BEFORE(tg_event)

Returns true if the trigger fired before the operation.

TRIGGER_FIRED_AFTER(tg_event)

Returns true if the trigger fired after the operation.

TRIGGER_FIRED_INSTEAD(tg_event)

Returns true if the trigger fired instead of the operation.

TRIGGER_FIRED_FOR_ROW(tg_event)

Returns true if the trigger fired for a row-level event.

TRIGGER_FIRED_FOR_STATEMENT(tg_event)

Returns true if the trigger fired for a statement-level event.

TRIGGER_FIRED_BY_INSERT(tg_event)

Returns true if the trigger was fired by an INSERT command.

TRIGGER_FIRED_BY_UPDATE(tg_event)

Returns true if the trigger was fired by an UPDATE command.

TRIGGER_FIRED_BY_DELETE(tg_event)

Returns true if the trigger was fired by a DELETE command.

TRIGGER_FIRED_BY_TRUNCATE(tg_event)

Returns true if the trigger was fired by a TRUNCATE command.

tg_relation

A pointer to a structure describing the relation that the trigger fired for. Look at utils/rel.h for details about this structure. The most interesting things are tg_relation->rd_att (de-
Triggers

scriptor of the relation tuples) and \( \text{tg}_\text{relation} \rightarrow \text{rd}_\text{rel} \rightarrow \text{relname} \) (relation name; the type is not \text{char*} but \text{NameData}; use \text{SPI_getrelname(tg\_relation)} to get a \text{char*} if you need a copy of the name).

\text{tg\_trigtuple}

A pointer to the row for which the trigger was fired. This is the row being inserted, updated, or deleted. If this trigger was fired for an \text{INSERT} or \text{DELETE} then this is what you should return from the function if you don’t want to replace the row with a different one (in the case of \text{INSERT}) or skip the operation. For triggers on foreign tables, values of system columns herein are unspecified.

\text{tg\_newtuple}

A pointer to the new version of the row, if the trigger was fired for an \text{UPDATE}, and \text{NULL} if it is for an \text{INSERT} or a \text{DELETE}. This is what you have to return from the function if the event is an \text{UPDATE} and you don’t want to replace this row by a different one or skip the operation. For triggers on foreign tables, values of system columns herein are unspecified.

\text{tg\_trigger}

A pointer to a structure of type \text{Trigger}, defined in \text{utils/reltrigger.h}:

\begin{verbatim}
typedef struct Trigger {
    Oid         tgoid;
    char       *tgname;
    Oid         tgfoid;
    int16       tgtype;
    char        tgenabled;
    bool       tgisinternal;
    Oid         tgconstrelid;
    Oid         tgconstindid;
    Oid         tgconstraint;
    bool        tgdeferrable;
    bool       tginitdeferred;
    int16       tgnargs;
    int16       tgnattr;
    int16       *tgattr;
    char        **tgargs;
    char        *tqqual;
    char        *tgoldtable;
    char        *tgnewtable;
} Trigger;
\end{verbatim}

where \text{tgname} is the trigger’s name, \text{tgnargs} is the number of arguments in \text{tg\_args}, and \text{tg\_args} is an array of pointers to the arguments specified in the \text{CREATE TRIGGER} statement. The other members are for internal use only.

\text{tg\_trigtuplebuf}

The buffer containing \text{tg\_trigtuple}, or \text{InvalidBuffer} if there is no such tuple or it is not stored in a disk buffer.

\text{tg\_newtuplebuf}

The buffer containing \text{tg\_newtuple}, or \text{InvalidBuffer} if there is no such tuple or it is not stored in a disk buffer.
A pointer to a structure of type `Tuplestorestate` containing zero or more rows in the format specified by `tg_relation`, or a NULL pointer if there is no OLD TABLE transition relation.

A pointer to a structure of type `Tuplestorestate` containing zero or more rows in the format specified by `tg_relation`, or a NULL pointer if there is no NEW TABLE transition relation.

To allow queries issued through SPI to reference transition tables, see SPI_register_trigger_data.

A trigger function must return either a `HeapTuple` pointer or a NULL pointer (not an SQL null value, that is, do not set `isNull` true). Be careful to return either `tg_trigtuple` or `tg_newtuple`, as appropriate, if you don't want to modify the row being operated on.

### 39.4. A Complete Trigger Example

Here is a very simple example of a trigger function written in C. (Examples of triggers written in procedural languages can be found in the documentation of the procedural languages.)

The function `trigf` reports the number of rows in the table `ttest` and skips the actual operation if the command attempts to insert a null value into the column `x`. (So the trigger acts as a not-null constraint but doesn't abort the transaction.)

First, the table definition:

```
CREATE TABLE ttest (
  x integer
);
```

This is the source code of the trigger function:

```c
#include "postgres.h"
#include "fmgr.h"
#include "executor/spi.h"       /* this is what you need to work
  with SPI */
#include "commands/trigger.h"   /* ... triggers ... */
#include "utils/rel.h"          /* ... and relations */

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(trigf);

Datum
trigf(PG_FUNCTION_ARGS)
{
  TriggerData *trigdata = (TriggerData *) fcinfo->context;
  TupleDesc   tupdesc;
  HeapTuple   rettuple;
  char        *when;
  bool        checknull = false;
  bool        isnull;
  int         ret, i;

  /* make sure it's called as a trigger at all */
  if (!CALLED_AS_TRIGGER(fcinfo))
```

1116

Triggers
elog(ERROR, "trigf: not called by trigger manager");

/* tuple to return to executor */
if (TRIGGER_FIRED_BY_UPDATE(trigdata->tg_event))
    rettuple = trigdata->tg_newtuple;
else
    rettuple = trigdata->tg_trigtuple;

/* check for null values */
if (!TRIGGER_FIRED_BY_DELETE(trigdata->tg_event)
    && TRIGGER_FIRED_BEFORE(trigdata->tg_event))
    checknull = true;

if (TRIGGER_FIRED_BEFORE(trigdata->tg_event))
    when = "before";
else
    when = "after ";

tupdesc = trigdata->tg_relation->rd_att;

/* connect to SPI manager */
if ((ret = SPI_connect()) < 0)
    elog(ERROR, "trigf (fired %s): SPI_connect returned %d", when, ret);

/* get number of rows in table */
ret = SPI_exec("SELECT count(*) FROM ttest", 0);

if (ret < 0)
    elog(ERROR, "trigf (fired %s): SPI_exec returned %d", when, ret);

/* count(*) returns int8, so be careful to convert */
i = DatumGetInt64(SPI_getbinval(SPI_tuptable->vals[0],
    SPI_tuptable->tupdesc, 1,
    &isnull));

elog (INFO, "trigf (fired %s): there are %d rows in ttest", when, i);

SPI_finish();

if (checknull)
{
    SPI_getbinval(rettuple, tupdesc, 1, &isnull);
    if (isnull)
        rettuple = NULL;
}

return PointerGetDatum(rettuple);

After you have compiled the source code (see Section 38.10.5), declare the function and the triggers:

CREATE FUNCTION trigf() RETURNS trigger
AS 'filename'
LANGUAGE C;

CREATE TRIGGER tbefore BEFORE INSERT OR UPDATE OR DELETE ON ttest
  FOR EACH ROW EXECUTE FUNCTION trigf();

CREATE TRIGGER tafter AFTER INSERT OR UPDATE OR DELETE ON ttest
  FOR EACH ROW EXECUTE FUNCTION trigf();

Now you can test the operation of the trigger:

=> INSERT INTO ttest VALUES (NULL);
INFO: trigf (fired before): there are 0 rows in ttest
INSERT 0 0
-- Insertion skipped and AFTER trigger is not fired

=> SELECT * FROM ttest;
 x
---
(0 rows)

=> INSERT INTO ttest VALUES (1);
INFO: trigf (fired before): there are 0 rows in ttest
INFO: trigf (fired after ): there are 1 rows in ttest

   remember what we said about

   visibility.
INSERT 167793 1
vac=> SELECT * FROM ttest;
 x
---
1
 (1 row)

=> INSERT INTO ttest SELECT x * 2 FROM ttest;
INFO: trigf (fired before): there are 1 rows in ttest
INFO: trigf (fired after ): there are 2 rows in ttest

   remember what we said about

   visibility.
INSERT 167794 1
=> SELECT * FROM ttest;
 x
---
1
 2
(2 rows)

=> UPDATE ttest SET x = NULL WHERE x = 2;
INFO: trigf (fired before): there are 2 rows in ttest
UPDATE 0
=> UPDATE ttest SET x = 4 WHERE x = 2;
INFO: trigf (fired before): there are 2 rows in ttest
INFO: trigf (fired after ): there are 2 rows in ttest
UPDATE 1
vac=> SELECT * FROM ttest;
Triggers

x
---
1
4
(2 rows)

=> DELETE FROM ttest;
INFO: trigf (fired before): there are 2 rows in ttest
INFO: trigf (fired before): there are 1 rows in ttest
INFO: trigf (fired after ): there are 0 rows in ttest
INFO: trigf (fired after ): there are 0 rows in ttest

^~~~

remember what we said about visibility.

DELETE 2
=> SELECT * FROM ttest;
  x
---
(0 rows)

There are more complex examples in src/test/regress/regress.c and in spi.
Chapter 40. Event Triggers

To supplement the trigger mechanism discussed in Chapter 39, PostgreSQL also provides event triggers. Unlike regular triggers, which are attached to a single table and capture only DML events, event triggers are global to a particular database and are capable of capturing DDL events.

Like regular triggers, event triggers can be written in any procedural language that includes event trigger support, or in C, but not in plain SQL.

40.1. Overview of Event Trigger Behavior

An event trigger fires whenever the event with which it is associated occurs in the database in which it is defined. Currently, the only supported events are `ddl_command_start`, `ddl_command_end`, `table_rewrite` and `sql_drop`. Support for additional events may be added in future releases.

The `ddl_command_start` event occurs just before the execution of a `CREATE`, `ALTER`, `DROP`, `SECURITY LABEL`, `COMMENT`, `GRANT` or `REVOKE` command. No check whether the affected object exists or doesn't exist is performed before the event trigger fires. As an exception, however, this event does not occur for DDL commands targeting shared objects — databases, roles, and tablespaces — or for commands targeting event triggers themselves. The event trigger mechanism does not support these object types. `ddl_command_start` also occurs just before the execution of a `SELECT INTO` command, since this is equivalent to `CREATE TABLE AS`.

The `ddl_command_end` event occurs just after the execution of this same set of commands. To obtain more details on the DDL operations that took place, use the set-returning function `pg_event_trigger_ddl_commands()` from the `ddl_command_end` event trigger code (see Section 9.28). Note that the trigger fires after the actions have taken place (but before the transaction commits), and thus the system catalogs can be read as already changed.

The `sql_drop` event occurs just before the `ddl_command_end` event trigger for any operation that drops database objects. To list the objects that have been dropped, use the set-returning function `pg_event_trigger_dropped_objects()` from the `sql_drop` event trigger code (see Section 9.28). Note that the trigger is executed after the objects have been deleted from the system catalogs, so it's not possible to look them up anymore.

The `table_rewrite` event occurs just before a table is rewritten by some actions of the commands `ALTER TABLE` and `ALTER TYPE`. While other control statements are available to rewrite a table, like `CLUSTER` and `VACUUM`, the `table_rewrite` event is not triggered by them.

Event triggers (like other functions) cannot be executed in an aborted transaction. Thus, if a DDL command fails with an error, any associated `ddl_command_end` triggers will not be executed. Conversely, if a `ddl_command_start` trigger fails with an error, no further event triggers will fire, and no attempt will be made to execute the command itself. Similarly, if a `ddl_command_end` trigger fails with an error, the effects of the DDL statement will be rolled back, just as they would be in any other case where the containing transaction aborts.

For a complete list of commands supported by the event trigger mechanism, see Section 40.2.

Event triggers are created using the command `CREATE EVENT TRIGGER`. In order to create an event trigger, you must first create a function with the special return type `event_trigger`. This function need not (and may not) return a value; the return type serves merely as a signal that the function is to be invoked as an event trigger.

If more than one event trigger is defined for a particular event, they will fire in alphabetical order by trigger name.

A trigger definition can also specify a `WHEN` condition so that, for example, a `ddl_command_start` trigger can be fired only for particular commands which the user wishes to intercept. A common use of such triggers is to restrict the range of DDL operations which users may perform.
## 40.2. Event Trigger Firing Matrix

Table 40.1 lists all commands for which event triggers are supported.

### Table 40.1. Event Trigger Support by Command Tag

<table>
<thead>
<tr>
<th>Command Tag</th>
<th>ddl_command_start</th>
<th>ddl_command_end</th>
<th>sql_drop</th>
<th>table_rewrite</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>ALTER AGGREGATE</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER COLLATION</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER CONVERSION</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER DOMAIN</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER EXTENSION</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER FOREIGN DATA WRAPPER</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER FOREIGN TABLE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER FUNCTION</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER LANGUAGE</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER OPERATOR</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER OPERATOR CLASS</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER OPERATOR FAMILY</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER POLICY</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER SCHEMA</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER SEQUENCE</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER SERVER</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>ALTER TABLE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>ALTER TEXT SEARCH</td>
<td>X</td>
<td>X</td>
<td>–</td>
<td>–</td>
<td></td>
</tr>
<tr>
<td>Command Tag</td>
<td>Ext: ddl_command_start</td>
<td>Ext: ddl_command_end</td>
<td>Ext: sql_drop</td>
<td>Ext: ta-ble_rewrite</td>
<td>Notes</td>
</tr>
<tr>
<td>-------------</td>
<td>------------------------</td>
<td>----------------------</td>
<td>--------------</td>
<td>--------------------</td>
<td>-------</td>
</tr>
<tr>
<td>CONFIGURATION</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER TEXT SEARCH DICTIONARY</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER TEXT SEARCH PARSER</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER TEXT SEARCH TEMPLATE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER TRIGGER</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER TYPE</td>
<td>X</td>
<td>X</td>
<td></td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>ALTER USER MAPPING</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER VIEW</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE AGGREGATE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>COMMENT</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td>Only for local objects</td>
</tr>
<tr>
<td>CREATE CAST</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE COLLATION</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE CONVERSION</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE DOMAIN</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE EXTENSION</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE FOREIGN DATA WRAPPER</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE FOREIGN TABLE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE FUNCTION</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE INDEX</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE LANGUAGE</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CREATE OPERATOR</td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Command Tag</td>
<td>ddl_command_start</td>
<td>ddl_command_end</td>
<td>sql_drop</td>
<td>table_rewrite</td>
<td>Notes</td>
</tr>
<tr>
<td>----------------------</td>
<td>-------------------</td>
<td>----------------</td>
<td>---------</td>
<td>---------------</td>
<td>-------</td>
</tr>
<tr>
<td>CREATE OPERATOR CLASS</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE OPERATOR FAMILY</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE POLICY</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE RULE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE SCHEMA</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE SEQUENCE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE SERVER</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE STATISTICS</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TABLE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TABLE AS</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TEXT SEARCH CONFIGURATION</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TEXT SEARCH DICTIONARY</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TEXT SEARCH PARSER</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TEMPLATE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TRIGGER</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE TYPE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE USER MAPPING</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>CREATE VIEW</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>Command Tag</td>
<td>ddl_command_start</td>
<td>ddl_command_end</td>
<td>sql_drop</td>
<td>table_rewrite</td>
<td>Notes</td>
</tr>
<tr>
<td>----------------------</td>
<td>-------------------</td>
<td>-----------------</td>
<td>----------</td>
<td>---------------</td>
<td>-------</td>
</tr>
<tr>
<td>DROP AGGREGATE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP CAST</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP DOMAIN</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP EXTENSION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP FOREIGN DATA WRAPPER</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP FOREIGN TABLE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP FUNCTION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP INDEX</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP LANGUAGE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP OPERATOR</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP OPERATOR CLASS</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP OPERATOR FAMILY</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP OWNED</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP POLICY</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP RULE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SCHEMA</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SEQUENCE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SERVER</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP STATISTICS</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TABLE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TEXT SEARCH CONFIGURATION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
</tbody>
</table>
### 40.3. Writing Event Trigger Functions in C

This section describes the low-level details of the interface to an event trigger function. This information is only needed when writing event trigger functions in C. If you are using a higher-level language then these details are handled for you. In most cases you should consider using a procedural language before writing your event triggers in C. The documentation of each procedural language explains how to write an event trigger in that language.

Event trigger functions must use the “version 1” function manager interface.

When a function is called by the event trigger manager, it is not passed any normal arguments, but it is passed a “context” pointer pointing to a `EventTriggerData` structure. C functions can check whether they were called from the event trigger manager or not by executing the macro:

```
CALLED_AS_EVENT_TRIGGER(fcinfo)
```

which expands to:

```
((fcinfo)->context != NULL && IsA((fcinfo)->context, EventTriggerData))
```
If this returns true, then it is safe to cast `fcinfo->context` to type `EventTriggerData` * and make use of the pointed-to `EventTriggerData` structure. The function must *not* alter the `EventTriggerData` structure or any of the data it points to.

```c
typedef struct EventTriggerData {
    NodeTag type;
    const char *event; /* event name */
    Node *parsetree; /* parse tree */
    const char *tag; /* command tag */
} EventTriggerData;
```

where the members are defined as follows:

**type**

Always T_EventTriggerData.

**event**

Describes the event for which the function is called, one of "ddl_command_start", "ddl_command_end", "sql_drop", "table_rewrite". See Section 40.1 for the meaning of these events.

**parsetree**

A pointer to the parse tree of the command. Check the PostgreSQL source code for details. The parse tree structure is subject to change without notice.

**tag**

The command tag associated with the event for which the event trigger is run, for example "CREATE FUNCTION".

An event trigger function must return a NULL pointer (*not* an SQL null value, that is, do not set `isNull` true).

### 40.4. A Complete Event Trigger Example

Here is a very simple example of an event trigger function written in C. (Examples of triggers written in procedural languages can be found in the documentation of the procedural languages.)

The function `noddl` raises an exception each time it is called. The event trigger definition associated the function with the `ddl_command_start` event. The effect is that all DDL commands (with the exceptions mentioned in Section 40.1) are prevented from running.

This is the source code of the trigger function:

```c
#include "postgres.h"
#include "commands/event_trigger.h"

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(noddl);
```
Datum

nodd1(PG_FUNCTION_ARGS)
{
    EventTriggerData *trigdata;

    if (!CALLED_AS_EVENT_TRIGGER(fcinfo)) /* internal error */
        elog(ERROR, "not fired by event trigger manager");

    trigdata = (EventTriggerData *) fcinfo->context;

    ereport(ERROR,
        (errcode(ERRCODE_INSUFFICIENT_PRIVILEGE),
        errmsg("command \"%s\" denied", trigdata->tag)));

    PG_RETURN_NULL();
}

After you have compiled the source code (see Section 38.10.5), declare the function and the triggers:

CREATE FUNCTION nodd1() RETURNS event_trigger
AS 'nodd1' LANGUAGE C;

CREATE EVENT TRIGGER nodd1 ON ddl_command_start
    EXECUTE FUNCTION nodd1();

Now you can test the operation of the trigger:

=# \dy
List of event triggers
Name       |       Event       | Owner | Enabled | Function | Tags
-------    |-------------------+-------+---------+----------+------
nodd1      | ddl_command_start | dim   | enabled | nodd1    |     
(1 row)

=# CREATE TABLE foo(id serial);
ERROR: command "CREATE TABLE" denied

In this situation, in order to be able to run some DDL commands when you need to do so, you have
to either drop the event trigger or disable it. It can be convenient to disable the trigger for only the
duration of a transaction:

BEGIN;
ALTER EVENT TRIGGER nodd1 DISABLE;
CREATE TABLE foo (id serial);
ALTER EVENT TRIGGER nodd1 ENABLE;
COMMIT;

(Recall that DDL commands on event triggers themselves are not affected by event triggers.)

40.5. A Table Rewrite Event Trigger Example

Thanks to the table_rewrite event, it is possible to implement a table rewriting policy only al-
lowing the rewrite in maintenance windows.

Here's an example implementing such a policy.
CREATE OR REPLACE FUNCTION no_rewrite()
RETURNS event_trigger
LANGUAGE plpgsql AS
$$
---
--- Implement local Table Rewriting policy:
--- public.foo is not allowed rewriting, ever
--- other tables are only allowed rewriting between 1am and 6am
--- unless they have more than 100 blocks
---
DECLARE
  table_oid oid := pg_event_trigger_table_rewrite_oid();
  current_hour integer := extract('hour' from current_time);
  pages integer;
  max_pages integer := 100;
BEGIN
  IF pg_event_trigger_table_rewrite_oid() = 'public.foo'::regclass
    THEN
      RAISE EXCEPTION 'you''re not allowed to rewrite the table
%'',
          table_oid::regclass;
    END IF;

  SELECT INTO pages relpages FROM pg_class WHERE oid = table_oid;
    IF pages > max_pages
      THEN
        RAISE EXCEPTION 'rewrites only allowed for table with less
than % pages'',
            max_pages;
      END IF;

  IF current_hour NOT BETWEEN 1 AND 6
    THEN
      RAISE EXCEPTION 'rewrites only allowed between 1am and
6am'';
    END IF;
END;
$$;

CREATE EVENT TRIGGER no_rewrite_allowed
  ON table_rewrite
  EXECUTE FUNCTION no_rewrite();
Chapter 41. The Rule System

This chapter discusses the rule system in PostgreSQL. Production rule systems are conceptually simple, but there are many subtle points involved in actually using them.

Some other database systems define active database rules, which are usually stored procedures and triggers. In PostgreSQL, these can be implemented using functions and triggers as well.

The rule system (more precisely speaking, the query rewrite rule system) is totally different from stored procedures and triggers. It modifies queries to take rules into consideration, and then passes the modified query to the query planner for planning and execution. It is very powerful, and can be used for many things such as query language procedures, views, and versions. The theoretical foundations and the power of this rule system are also discussed in [ston90b] and [ong90].

41.1. The Query Tree

To understand how the rule system works it is necessary to know when it is invoked and what its input and results are.

The rule system is located between the parser and the planner. It takes the output of the parser, one query tree, and the user-defined rewrite rules, which are also query trees with some extra information, and creates zero or more query trees as result. So its input and output are always things the parser itself could have produced and thus, anything it sees is basically representable as an SQL statement.

Now what is a query tree? It is an internal representation of an SQL statement where the single parts that it is built from are stored separately. These query trees can be shown in the server log if you set the configuration parameters debug_print_parse, debug_print_rewritten, or debug_print_plan. The rule actions are also stored as query trees, in the system catalog pg_rewrite. They are not formatted like the log output, but they contain exactly the same information.

Reading a raw query tree requires some experience. But since SQL representations of query trees are sufficient to understand the rule system, this chapter will not teach how to read them.

When reading the SQL representations of the query trees in this chapter it is necessary to be able to identify the parts the statement is broken into when it is in the query tree structure. The parts of a query tree are

the command type

This is a simple value telling which command (SELECT, INSERT, UPDATE, DELETE) produced the query tree.

the range table

The range table is a list of relations that are used in the query. In a SELECT statement these are the relations given after the FROM key word.

Every range table entry identifies a table or view and tells by which name it is called in the other parts of the query. In the query tree, the range table entries are referenced by number rather than by name, so here it doesn't matter if there are duplicate names as it would in an SQL statement. This can happen after the range tables of rules have been merged in. The examples in this chapter will not have this situation.

the result relation

This is an index into the range table that identifies the relation where the results of the query go.

SELECT queries don't have a result relation. (The special case of SELECT INTO is mostly identical to CREATE TABLE followed by INSERT ... SELECT, and is not discussed separately here.)
For INSERT, UPDATE, and DELETE commands, the result relation is the table (or view!) where the changes are to take effect.

the target list

The target list is a list of expressions that define the result of the query. In the case of a SELECT, these expressions are the ones that build the final output of the query. They correspond to the expressions between the key words SELECT and FROM. (* is just an abbreviation for all the column names of a relation. It is expanded by the parser into the individual columns, so the rule system never sees it.)

DELETE commands don't need a normal target list because they don't produce any result. Instead, the planner adds a special CTID entry to the empty target list, to allow the executor to find the row to be deleted. (CTID is added when the result relation is an ordinary table. If it is a view, a whole-row variable is added instead, by the rule system, as described in Section 41.2.4.)

For INSERT commands, the target list describes the new rows that should go into the result relation. It consists of the expressions in the VALUES clause or the ones from the SELECT clause in INSERT ... SELECT. The first step of the rewrite process adds target list entries for any columns that were not assigned to by the original command but have defaults. Any remaining columns (with neither a given value nor a default) will be filled in by the planner with a constant null expression.

For UPDATE commands, the target list describes the new rows that should replace the old ones. In the rule system, it contains just the expressions from the SET column = expression part of the command. The planner will handle missing columns by inserting expressions that copy the values from the old row into the new one. Just as for DELETE, a CTID or whole-row variable is added so that the executor can identify the old row to be updated.

Every entry in the target list contains an expression that can be a constant value, a variable pointing to a column of one of the relations in the range table, a parameter, or an expression tree made of function calls, constants, variables, operators, etc.

the qualification

The query's qualification is an expression much like one of those contained in the target list entries. The result value of this expression is a Boolean that tells whether the operation (INSERT, UPDATE, DELETE, or SELECT) for the final result row should be executed or not. It corresponds to the WHERE clause of an SQL statement.

the join tree

The query's join tree shows the structure of the FROM clause. For a simple query like SELECT ... FROM a, b, c, the join tree is just a list of the FROM items, because we are allowed to join them in any order. But when JOIN expressions, particularly outer joins, are used, we have to join in the order shown by the joins. In that case, the join tree shows the structure of the JOIN expressions. The restrictions associated with particular JOIN clauses (from ON or USING expressions) are stored as qualification expressions attached to those join-tree nodes. It turns out to be convenient to store the top-level WHERE expression as a qualification attached to the top-level join-tree item, too. So really the join tree represents both the FROM and WHERE clauses of a SELECT.

the others

The other parts of the query tree like the ORDER BY clause aren't of interest here. The rule system substitutes some entries there while applying rules, but that doesn't have much to do with the fundamentals of the rule system.

41.2. Views and the Rule System
Views in PostgreSQL are implemented using the rule system. In fact, there is essentially no difference between:

```sql
CREATE VIEW myview AS SELECT * FROM mytab;
```

compared against the two commands:

```sql
CREATE TABLE myview (same column list as mytab);
CREATE RULE "_RETURN" AS ON SELECT TO myview DO INSTEAD
   SELECT * FROM mytab;
```

because this is exactly what the CREATE VIEW command does internally. This has some side effects. One of them is that the information about a view in the PostgreSQL system catalogs is exactly the same as it is for a table. So for the parser, there is absolutely no difference between a table and a view. They are the same thing: relations.

### 41.2.1. How SELECT Rules Work

Rules ON SELECT are applied to all queries as the last step, even if the command given is an INSERT, UPDATE or DELETE. And they have different semantics from rules on the other command types in that they modify the query tree in place instead of creating a new one. So SELECT rules are described first.

Currently, there can be only one action in an ON SELECT rule, and it must be an unconditional SELECT action that is INSTEAD. This restriction was required to make rules safe enough to open them for ordinary users, and it restricts ON SELECT rules to act like views.

The examples for this chapter are two join views that do some calculations and some more views using them in turn. One of the two first views is customized later by adding rules for INSERT, UPDATE, and DELETE operations so that the final result will be a view that behaves like a real table with some magic functionality. This is not such a simple example to start from and this makes things harder to get into. But it's better to have one example that covers all the points discussed step by step rather than having many different ones that might mix up in mind.

For the example, we need a little `min` function that returns the lower of 2 integer values. We create that as:

```sql
CREATE FUNCTION min(integer, integer) RETURNS integer AS $$
   SELECT CASE WHEN $1 < $2 THEN $1 ELSE $2 END
$$ LANGUAGE SQL STRICT;
```

The real tables we need in the first two rule system descriptions are these:

```sql
CREATE TABLE shoe_data (  
   shoename text,          -- primary key  
   sh_avail integer,       -- available number of pairs  
   slcolor text,           -- preferred shoelace color  
   slminlen real,          -- minimum shoelace length  
   slmaxlen real,          -- maximum shoelace length  
   slunit text             -- length unit
);
```

```sql
CREATE TABLE shoelace_data (  
   sl_name    text,          -- primary key  
   sl_avail   integer,       -- available number of pairs  
   sl_color   text,          -- shoelace color  
   sl_len     real,          -- shoelace length  
   sl_unit    text           -- length unit
);  
```
The Rule System

CREATE TABLE unit (  
    un_name    text,          -- primary key  
    un_fact    real           -- factor to transform to cm  
);  

As you can see, they represent shoe-store data.

The views are created as:

CREATE VIEW shoe AS  
SELECT sh.shoename,  
    sh.sh_avail,  
    sh.slcolor,  
    sh.slminlen,  
    sh.slminlen * un.un_fact AS slminlen_cm,  
    sh.slmaxlen,  
    sh.slmaxlen * un.un_fact AS slmaxlen_cm,  
    sh.slunit  
FROM shoe_data sh, unit un  
WHERE sh.slunit = un.un_name;

CREATE VIEW shoelace AS  
SELECT s.sl_name,  
    s.sl_avail,  
    s.sl_color,  
    s.sl_len,  
    s.sl_unit,  
    s.sl_len * u.un_fact AS sl_len_cm  
FROM shoelace_data s, unit u  
WHERE s.sl_unit = u.un_name;

CREATE VIEW shoe_ready AS  
SELECT rsh.shoename,  
    rsh.sh_avail,  
    rsl.sl_name,  
    rsl.sl_avail,  
    min(rsh.sh_avail, rsl.sl_avail) AS total_avail  
FROM shoe rsh, shoelace rsl  
WHERE rsl.sl_color = rsh.slcolor  
AND rsl.sl_len_cm >= rsh.slminlen_cm  
AND rsl.sl_len_cm <= rsh.slmaxlen_cm;

The CREATE VIEW command for the shoelace view (which is the simplest one we have) will create a relation shoelace and an entry in pg_rewrite that tells that there is a rewrite rule that must be applied whenever the relation shoelace is referenced in a query’s range table. The rule has no rule qualification (discussed later, with the non-SELECT rules, since SELECT rules currently cannot have them) and it is INSTEAD. Note that rule qualifications are not the same as query qualifications. The action of our rule has a query qualification. The action of the rule is one query tree that is a copy of the SELECT statement in the view creation command.

Note

The two extra range table entries for NEW and OLD that you can see in the pg_rewrite entry aren’t of interest for SELECT rules.
Now we populate `unit`, `shoe_data` and `shoelace_data` and run a simple query on a view:

```sql
INSERT INTO unit VALUES ('cm', 1.0);
INSERT INTO unit VALUES ('m', 100.0);
INSERT INTO unit VALUES ('inch', 2.54);

INSERT INTO shoe_data VALUES ('sh1', 2, 'black', 70.0, 90.0, 'cm');
INSERT INTO shoe_data VALUES ('sh2', 0, 'black', 30.0, 40.0, 'inch');
INSERT INTO shoe_data VALUES ('sh3', 4, 'brown', 50.0, 65.0, 'cm');
INSERT INTO shoe_data VALUES ('sh4', 3, 'brown', 40.0, 50.0, 'inch');

INSERT INTO shoelace_data VALUES ('sl1', 5, 'black', 80.0, 'cm');
INSERT INTO shoelace_data VALUES ('sl2', 6, 'black', 100.0, 'cm');
INSERT INTO shoelace_data VALUES ('sl3', 0, 'black', 35.0, 'inch');
INSERT INTO shoelace_data VALUES ('sl4', 8, 'black', 40.0, 'inch');
INSERT INTO shoelace_data VALUES ('sl5', 4, 'brown', 1.0, 'm');
INSERT INTO shoelace_data VALUES ('sl6', 0, 'brown', 0.9, 'm');
INSERT INTO shoelace_data VALUES ('sl7', 7, 'brown', 60, 'cm');
INSERT INTO shoelace_data VALUES ('sl8', 1, 'brown', 40, 'inch');

SELECT * FROM shoelace;
```

```
<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl7</td>
<td>7</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl3</td>
<td>0</td>
<td>black</td>
<td>35</td>
<td>inch</td>
<td>88.9</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl8</td>
<td>1</td>
<td>brown</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl15</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl16</td>
<td>0</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>
```

This is the simplest `SELECT` you can do on our views, so we take this opportunity to explain the basics of view rules. The `SELECT * FROM shoelace` was interpreted by the parser and produced the query tree:

```sql
SELECT shoelace.sl_name, shoelace.sl_avail, shoelace.sl_color, shoelace.sl_len, shoelace.sl_unit, shoelace.sl_len_cm FROM shoelace shoelace;
```

and this is given to the rule system. The rule system walks through the range table and checks if there are rules for any relation. When processing the range table entry for `shoelace` (the only one up to now) it finds the _RETURN rule with the query tree:

```sql
SELECT s.sl_name, s.sl_avail, s.sl_color, s.sl_len, s.sl_unit, s.sl_len * u.un_fact AS sl_len_cm FROM shoelace old, shoelace new, shoelace_data s, unit u
```

1133
WHERE s.sl_unit = u.un_name;

To expand the view, the rewriter simply creates a subquery range-table entry containing the rule’s action query tree, and substitutes this range table entry for the original one that referenced the view. The resulting rewritten query tree is almost the same as if you had typed:

```
SELECT shoelace.sl_name, shoelace.sl_avail,
      shoelace.sl_color, shoelace.sl_len,
      shoelace.sl_unit, shoelace.sl_len_cm
FROM (SELECT s.sl_name,
      s.sl_avail,
      s.sl_color,
      s.sl_len,
      s.sl_unit,
      s.sl_len * u.un_fact AS sl_len_cm
      FROM shoelace_data s, unit u
      WHERE s.sl_unit = u.un_name) shoelace;
```

There is one difference however: the subquery’s range table has two extra entries shoelace old and shoelace new. These entries don’t participate directly in the query, since they aren’t referenced by the subquery’s join tree or target list. The rewriter uses them to store the access privilege check information that was originally present in the range-table entry that referenced the view. In this way, the executor will still check that the user has proper privileges to access the view, even though there’s no direct use of the view in the rewritten query.

That was the first rule applied. The rule system will continue checking the remaining range-table entries in the top query (in this example there are no more), and it will recursively check the range-table entries in the added subquery to see if any of them reference views. (But it won’t expand old or new — otherwise we’d have infinite recursion!) In this example, there are no rewrite rules for shoelace_data or unit, so rewriting is complete and the above is the final result given to the planner.

Now we want to write a query that finds out for which shoes currently in the store we have the matching shoelaces (color and length) and where the total number of exactly matching pairs is greater or equal to two.

```
SELECT * FROM shoe_ready WHERE total_avail >= 2;
```

```
shoename | sh_avail | sl_name | sl_avail | total_avail
----------+----------+---------+----------+-------------
sh1      |        2 | sl1     |        5 |           2
sh3      |        4 | sl7     |        7 |           4
(2 rows)
```

The output of the parser this time is the query tree:

```
SELECT shoe_ready.shoename, shoe_ready.sh_avail,
      shoe_ready.sl_name, shoe_ready.sl_avail,
      shoe_ready.total_avail
FROM shoe_ready shoe_ready
WHERE shoe_ready.total_avail >= 2;
```

The first rule applied will be the one for the shoe_ready view and it results in the query tree:

```
SELECT shoe_ready.shoename, shoe_ready.sh_avail,
      shoe_ready.sl_name, shoe_ready.sl_avail,
```
The Rule System

shoe_ready.total_avail
FROM (SELECT rsh.shoename,
rsh.sh_avail,
rsl.sl_name,
rsl.sl_avail,
min(rsh.sh_avail, rsl.sl_avail) AS total_avail
FROM shoe rsh, shoelace rsl
WHERE rsl.sl_color = rsh.slcolor
AND rsl.sl_len_cm >= rsh.slminlen_cm
AND rsl.sl_len_cm <= rsh.slmaxlen_cm) shoe_ready
WHERE shoe_ready.total_avail >= 2;

Similarly, the rules for shoe and shoelace are substituted into the range table of the subquery, leading to a three-level final query tree:

SELECT shoe_ready.shoename, shoe_ready.sh_avail,
shoe_ready.sl_name, shoe_ready.sl_avail,
shoe_ready.total_avail
FROM (SELECT rsh.shoename,
rsh.sh_avail,
rsl.sl_name,
rsl.sl_avail,
min(rsh.sh_avail, rsl.sl_avail) AS total_avail
FROM (SELECT sh.shoename,
sh.sh_avail,
sh.slcolor,
sh.slminlen,
sh.slminlen * un.un_fact AS slminlen_cm,
sh.slmaxlen,
sh.slmaxlen * un.un_fact AS slmaxlen_cm,
sh.slunit
FROM shoe_data sh, unit un
WHERE sh.slunit = un.un_name) rsh,
(SELECT s.sl_name,
s.sl_avail,
s.sl_color,
s.sl_len,
s.sl_unit,
s.sl_len * u.un_fact AS sl_len_cm
FROM shoelace_data s, unit u
WHERE s.sl_unit = u.un_name) rsl
WHERE rsl.sl_color = rsh.slcolor
AND rsl.sl_len_cm >= rsh.slminlen_cm
AND rsl.sl_len_cm <= rsh.slmaxlen_cm) shoe_ready
WHERE shoe_ready.total_avail >= 2;

It turns out that the planner will collapse this tree into a two-level query tree: the bottommost SELECT commands will be "pulled up" into the middle SELECT since there's no need to process them separately. But the middle SELECT will remain separate from the top, because it contains aggregate functions. If we pulled those up it would change the behavior of the topmost SELECT, which we don't want. However, collapsing the query tree is an optimization that the rewrite system doesn't have to concern itself with.

41.2.2. View Rules in Non-SELECT Statements

Two details of the query tree aren't touched in the description of view rules above. These are the command type and the result relation. In fact, the command type is not needed by view rules, but the
result relation may affect the way in which the query rewriter works, because special care needs to be taken if the result relation is a view.

There are only a few differences between a query tree for a \texttt{SELECT} and one for any other command. Obviously, they have a different command type and for a command other than a \texttt{SELECT}, the result relation points to the range-table entry where the result should go. Everything else is absolutely the same. So having two tables \texttt{t1} and \texttt{t2} with columns \texttt{a} and \texttt{b}, the query trees for the two statements:

\begin{verbatim}
SELECT t2.b FROM t1, t2 WHERE t1.a = t2.a;
UPDATE t1 SET b = t2.b FROM t2 WHERE t1.a = t2.a;
\end{verbatim}

are nearly identical. In particular:

\begin{itemize}
\item The range tables contain entries for the tables \texttt{t1} and \texttt{t2}.
\item The target lists contain one variable that points to column \texttt{b} of the range table entry for table \texttt{t2}.
\item The qualification expressions compare the columns \texttt{a} of both range-table entries for equality.
\item The join trees show a simple join between \texttt{t1} and \texttt{t2}.
\end{itemize}

The consequence is, that both query trees result in similar execution plans: They are both joins over the two tables. For the \texttt{UPDATE} the missing columns from \texttt{t1} are added to the target list by the planner and the final query tree will read as:

\begin{verbatim}
UPDATE t1 SET a = t1.a, b = t2.b FROM t2 WHERE t1.a = t2.a;
\end{verbatim}

and thus the executor run over the join will produce exactly the same result set as:

\begin{verbatim}
SELECT t1.a, t2.b FROM t1, t2 WHERE t1.a = t2.a;
\end{verbatim}

But there is a little problem in \texttt{UPDATE}: the part of the executor plan that does the join does not care what the results from the join are meant for. It just produces a result set of rows. The fact that one is a \texttt{SELECT} command and the other is an \texttt{UPDATE} is handled higher up in the executor, where it knows that this is an \texttt{UPDATE}, and it knows that this result should go into table \texttt{t1}. But which of the rows that are there has to be replaced by the new row?

To resolve this problem, another entry is added to the target list in \texttt{UPDATE} (and also in \texttt{DELETE}) statements: the current tuple ID (CTID). This is a system column containing the file block number and position in the block for the row. Knowing the table, the CTID can be used to retrieve the original row of \texttt{t1} to be updated. After adding the CTID to the target list, the query actually looks like:

\begin{verbatim}
SELECT t1.a, t2.b, t1.ctid FROM t1, t2 WHERE t1.a = t2.a;
\end{verbatim}

Now another detail of PostgreSQL enters the stage. Old table rows aren't overwritten, and this is why \texttt{ROLLBACK} is fast. In an \texttt{UPDATE}, the new result row is inserted into the table (after stripping the CTID) and in the row header of the old row, which the CTID pointed to, the \texttt{cmax} and \texttt{xmax} entries are set to the current command counter and current transaction ID. Thus the old row is hidden, and after the transaction commits the vacuum cleaner can eventually remove the dead row.

Knowing all that, we can simply apply view rules in absolutely the same way to any command. There is no difference.

### 41.2.3. The Power of Views in PostgreSQL

The above demonstrates how the rule system incorporates view definitions into the original query tree. In the second example, a simple \texttt{SELECT} from one view created a final query tree that is a join of 4 tables (\texttt{unit} was used twice with different names).
The benefit of implementing views with the rule system is, that the planner has all the information about which tables have to be scanned plus the relationships between these tables plus the restrictive qualifications from the views plus the qualifications from the original query in one single query tree. And this is still the situation when the original query is already a join over views. The planner has to decide which is the best path to execute the query, and the more information the planner has, the better this decision can be. And the rule system as implemented in PostgreSQL ensures, that this is all information available about the query up to that point.

### 41.2.4. Updating a View

What happens if a view is named as the target relation for an `INSERT`, `UPDATE`, or `DELETE`? Doing the substitutions described above would give a query tree in which the result relation points at a subquery range-table entry, which will not work. There are several ways in which PostgreSQL can support the appearance of updating a view, however.

If the subquery selects from a single base relation and is simple enough, the rewriter can automatically replace the subquery with the underlying base relation so that the `INSERT`, `UPDATE`, or `DELETE` is applied to the base relation in the appropriate way. Views that are “simple enough” for this are called *automatically updatable*. For detailed information on the kinds of view that can be automatically updated, see `CREATE VIEW`.

Alternatively, the operation may be handled by a user-provided `INSTEAD OF` trigger on the view. Rewriting works slightly differently in this case. For `INSERT`, the rewriter does nothing at all with the view, leaving it as the result relation for the query. For `UPDATE` and `DELETE`, it’s still necessary to expand the view query to produce the “old” rows that the command will attempt to update or delete. So the view is expanded as normal, but another unexpanded range-table entry is added to the query to represent the view in its capacity as the result relation.

The problem that now arises is how to identify the rows to be updated in the view. Recall that when the result relation is a table, a special CTID entry is added to the target list to identify the physical locations of the rows to be updated. This does not work if the result relation is a view, because a view does not have any CTID, since its rows do not have actual physical locations. Instead, for an `UPDATE` or `DELETE` operation, a special `wholerow` entry is added to the target list, which expands to include all columns from the view. The executor uses this value to supply the “old” row to the `INSTEAD OF` trigger. It is up to the trigger to work out what to update based on the old and new row values.

Another possibility is for the user to define `INSTEAD` rules that specify substitute actions for `INSERT`, `UPDATE`, and `DELETE` commands on a view. These rules will rewrite the command, typically into a command that updates one or more tables, rather than views. That is the topic of Section 41.4.

Note that rules are evaluated first, rewriting the original query before it is planned and executed. Therefore, if a view has `INSTEAD OF` triggers as well as rules on `INSERT`, `UPDATE`, or `DELETE`, then the rules will be evaluated first, and depending on the result, the triggers may not be used at all.

Automatic rewriting of an `INSERT`, `UPDATE`, or `DELETE` query on a simple view is always tried last. Therefore, if a view has rules or triggers, they will override the default behavior of automatically updatable views.

If there are no `INSTEAD` rules or `INSTEAD OF` triggers for the view, and the rewriter cannot automatically rewrite the query as an update on the underlying base relation, an error will be thrown because the executor cannot update a view as such.

### 41.3. Materialized Views

Materialized views in PostgreSQL use the rule system like views do, but persist the results in a table-like form. The main differences between:

```sql
CREATE MATERIALIZED VIEW mymatview AS SELECT * FROM mytab;
```
and:

```sql
CREATE TABLE mymatview AS SELECT * FROM mytab;
```

are that the materialized view cannot subsequently be directly updated and that the query used to create
the materialized view is stored in exactly the same way that a view's query is stored, so that fresh data
can be generated for the materialized view with:

```sql
REFRESH MATERIALIZED VIEW mymatview;
```

The information about a materialized view in the PostgreSQL system catalogs is exactly the same as
it is for a table or view. So for the parser, a materialized view is a relation, just like a table or a view.
When a materialized view is referenced in a query, the data is returned directly from the materialized
view, like from a table; the rule is only used for populating the materialized view.

While access to the data stored in a materialized view is often much faster than accessing the underlying
tables directly or through a view, the data is not always current; yet sometimes current data is not
needed. Consider a table which records sales:

```sql
CREATE TABLE invoice (
    invoice_no    integer        PRIMARY KEY,
    seller_no     integer,       -- ID of salesperson
    invoice_date  date,          -- date of sale
    invoice_amt   numeric(13,2)  -- amount of sale
);
```

If people want to be able to quickly graph historical sales data, they might want to summarize, and
they may not care about the incomplete data for the current date:

```sql
CREATE MATERIALIZED VIEW sales_summary AS
    SELECT
        seller_no,
        invoice_date,
        sum(invoice_amt)::numeric(13,2) as sales_amt
    FROM invoice
    WHERE invoice_date < CURRENT_DATE
    GROUP BY
        seller_no,
        invoice_date
    ORDER BY
        seller_no,
        invoice_date;
```

This materialized view might be useful for displaying a graph in the dashboard created for salespeople.
A job could be scheduled to update the statistics each night using this SQL statement:

```sql
REFRESH MATERIALIZED VIEW sales_summary;
```

Another use for a materialized view is to allow faster access to data brought across from a remote
system through a foreign data wrapper. A simple example using file_fdw is below, with timings,
but since this is using cache on the local system the performance difference compared to access to a
remote system would usually be greater than shown here. Notice we are also exploiting the ability to
put an index on the materialized view, whereas file_fdw does not support indexes; this advantage might not apply for other sorts of foreign data access.

Setup:

```
CREATE EXTENSION file_fdw;
CREATE SERVER local_file FOREIGN DATA WRAPPER file_fdw;
CREATE FOREIGN TABLE words (word text NOT NULL)
    SERVER local_file
    OPTIONS (filename '/usr/share/dict/words');
CREATE MATERIALIZED VIEW wrd AS SELECT * FROM words;
CREATE UNIQUE INDEX wrd_word ON wrd (word);
CREATE EXTENSION pg_trgm;
CREATE INDEX wrd_trgm ON wrd USING gist (word gist_trgm_ops);
VACUUM ANALYZE wrd;
```

Now let’s spell-check a word. Using file_fdw directly:

```
SELECT count(*) FROM words WHERE word = 'caterpiler';

<table>
<thead>
<tr>
<th>count</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
</tr>
</tbody>
</table>

With EXPLAIN ANALYZE, we see:

```
Aggregate (cost=21763.99..21764.00 rows=1 width=0) (actual time=188.180..188.181 rows=1 loops=1)
  ->  Foreign Scan on words (cost=0.00..21761.41 rows=1032 width=0) (actual time=188.177..188.177 rows=0 loops=1)
        Filter: (word = 'caterpiler'::text)
        Rows Removed by Filter: 479829
        Foreign File: /usr/share/dict/words
        Foreign File Size: 4953699
Planning time: 0.118 ms
Execution time: 188.273 ms
```

If the materialized view is used instead, the query is much faster:

```
Aggregate (cost=4.44..4.45 rows=1 width=0) (actual time=0.042..0.042 rows=1 loops=1)
  ->  Index Only Scan using wrd_word on wrd (cost=0.42..4.44 rows=1 width=0) (actual time=0.039..0.039 rows=0 loops=1)
        Index Cond: (word = 'caterpiler'::text)
        Heap Fetches: 0
Planning time: 0.164 ms
Execution time: 0.117 ms
```

Either way, the word is spelled wrong, so let’s look for what we might have wanted. Again using file_fdw:

```
SELECT word FROM words ORDER BY word <-> 'caterpiler' LIMIT 10;
```

<table>
<thead>
<tr>
<th>word</th>
</tr>
</thead>
<tbody>
<tr>
<td>1139</td>
</tr>
</tbody>
</table>
41.4. Rules on INSERT, UPDATE, and DELETE

Rules that are defined on INSERT, UPDATE, and DELETE are significantly different from the view rules described in the previous section. First, their CREATE RULE command allows more:

• They are allowed to have no action.
• They can have multiple actions.
• They can be INSTEAD or ALSO (the default).
• The pseudorelations NEW and OLD become useful.
• They can have rule qualifications.

Second, they don’t modify the query tree in place. Instead they create zero or more new query trees and can throw away the original one.
Caution

In many cases, tasks that could be performed by rules on `INSERT/UPDATE/DELETE` are better done with triggers. Triggers are notationally a bit more complicated, but their semantics are much simpler to understand. Rules tend to have surprising results when the original query contains volatile functions: volatile functions may get executed more times than expected in the process of carrying out the rules.

Also, there are some cases that are not supported by these types of rules at all, notably including `WITH` clauses in the original query and multiple-assignment sub-`SELECT`s in the `SET` list of `UPDATE` queries. This is because copying these constructs into a rule query would result in multiple evaluations of the sub-query, contrary to the express intent of the query's author.

41.4.1. How Update Rules Work

Keep the syntax:

```
CREATE [ OR REPLACE ] RULE name AS ON event
    TO table [ WHERE condition ]
    DO [ ALSO | INSTEAD ] { NOTHING | command | ( command ; command ... ) }
```

in mind. In the following, `update rules` means rules that are defined on `INSERT, UPDATE, or DELETE`.

Update rules get applied by the rule system when the result relation and the command type of a query tree are equal to the object and event given in the `CREATE RULE` command. For update rules, the rule system creates a list of query trees. Initially the query-tree list is empty. There can be zero (NOTHING key word), one, or multiple actions. To simplify, we will look at a rule with one action. This rule can have a qualification or not and it can be `INSTEAD` or `ALSO` (the default).

What is a rule qualification? It is a restriction that tells when the actions of the rule should be done and when not. This qualification can only reference the pseudorelations `NEW` and/or `OLD`, which basically represent the relation that was given as object (but with a special meaning).

So we have three cases that produce the following query trees for a one-action rule.

- **No qualification, with either `ALSO` or `INSTEAD`**
  - the query tree from the rule action with the original query tree's qualification added

- **Qualification given and `ALSO`**
  - the query tree from the rule action with the rule qualification and the original query tree's qualification added

- **Qualification given and `INSTEAD`**
  - the query tree from the rule action with the rule qualification and the original query tree's qualification; and the original query tree with the negated rule qualification added

Finally, if the rule is `ALSO`, the unchanged original query tree is added to the list. Since only qualified `INSTEAD` rules already add the original query tree, we end up with either one or two output query trees for a rule with one action.

For `ON INSERT` rules, the original query (if not suppressed by `INSTEAD`) is done before any actions added by rules. This allows the actions to see the inserted row(s). But for `ON UPDATE` and `ON`
DELETE rules, the original query is done after the actions added by rules. This ensures that the actions can see the to-be-updated or to-be-deleted rows; otherwise, the actions might do nothing because they find no rows matching their qualifications.

The query trees generated from rule actions are thrown into the rewrite system again, and maybe more rules get applied resulting in more or less query trees. So a rule's actions must have either a different command type or a different result relation than the rule itself is on, otherwise this recursive process will end up in an infinite loop. (Recursive expansion of a rule will be detected and reported as an error.)

The query trees found in the actions of the pg_rewrite system catalog are only templates. Since they can reference the range-table entries for NEW and OLD, some substitutions have to be made before they can be used. For any reference to NEW, the target list of the original query is searched for a corresponding entry. If found, that entry's expression replaces the reference. Otherwise, NEW means the same as OLD (for an UPDATE) or is replaced by a null value (for an INSERT). Any reference to OLD is replaced by a reference to the range-table entry that is the result relation.

After the system is done applying update rules, it applies view rules to the produced query tree(s). Views cannot insert new update actions so there is no need to apply update rules to the output of view rewriting.

### 41.4.1.1. A First Rule Step by Step

Say we want to trace changes to the `sl_avail` column in the `shoelace_data` relation. So we set up a log table and a rule that conditionally writes a log entry when an UPDATE is performed on `shoelace_data`.

```sql
CREATE TABLE shoelace_log (
    sl_name    text,          -- shoelace changed
    sl_avail   integer,       -- new available value
    log_who    text,          -- who did it
    log_when   timestamp      -- when
);

CREATE RULE log_shoelace AS ON UPDATE TO shoelace_data
    WHERE NEW.sl_avail <> OLD.sl_avail
    DO INSERT INTO shoelace_log VALUES (
        NEW.sl_name,
        NEW.sl_avail,
        current_user,
        current_timestamp
    );
```

Now someone does:

```sql
UPDATE shoelace_data SET sl_avail = 6 WHERE sl_name = 'sl7';
```

and we look at the log table:

```sql
SELECT * FROM shoelace_log;
```

```plaintext
  sl_name | sl_avail | log_who | log_when
---------+----------+---------+----------------------------------
  sl7     |        6 | Al      | Tue Oct 20 16:14:45 1998 MET DST
```

That's what we expected. What happened in the background is the following. The parser created the query tree:
UPDATE shoelace_data SET sl_avail = 6
FROM shoelace_data shoelace_data
WHERE shoelace_data.sl_name = 'sl7';

There is a rule log_shoelace that is ON UPDATE with the rule qualification expression:

NEW.sl_avail <> OLD.sl_avail

and the action:

INSERT INTO shoelace_log VALUES (
    new.sl_name, new.sl_avail,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old;

(This looks a little strange since you cannot normally write INSERT ... VALUES ... FROM. The FROM clause here is just to indicate that there are range-table entries in the query tree for new and old. These are needed so that they can be referenced by variables in the INSERT command's query tree.)

The rule is a qualified ALSO rule, so the rule system has to return two query trees: the modified rule action and the original query tree. In step 1, the range table of the original query is incorporated into the rule’s action query tree. This results in:

INSERT INTO shoelace_log VALUES (
    new.sl_name, new.sl_avail,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
    shoelace_data shoelace_data;

In step 2, the rule qualification is added to it, so the result set is restricted to rows where sl_avail changes:

INSERT INTO shoelace_log VALUES (
    new.sl_name, new.sl_avail,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
    shoelace_data shoelace_data
WHERE new.sl_avail <> old.sl_avail;

(This looks even stranger, since INSERT ... VALUES doesn’t have a WHERE clause either, but the planner and executor will have no difficulty with it. They need to support this same functionality anyway for INSERT ... SELECT.)

In step 3, the original query tree’s qualification is added, restricting the result set further to only the rows that would have been touched by the original query:

INSERT INTO shoelace_log VALUES (
    new.sl_name, new.sl_avail,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
    shoelace_data shoelace_data
WHERE new.sl_avail <> old.sl_avail
    AND shoelace_data.sl_name = 'sl7';
Step 4 replaces references to NEW by the target list entries from the original query tree or by the matching variable references from the result relation:

```
INSERT INTO shoelace_log VALUES (
    shoelace_data.sl_name, 6,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
    shoelace_data shoelace_data
WHERE 6 <> old.sl_avail
    AND shoelace_data.sl_name = 'sl7';
```

Step 5 changes OLD references into result relation references:

```
INSERT INTO shoelace_log VALUES (
    shoelace_data.sl_name, 6,
    current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
    shoelace_data shoelace_data
WHERE 6 <> shoelace_data.sl_avail
    AND shoelace_data.sl_name = 'sl7';
```

That’s it. Since the rule is ALSO, we also output the original query tree. In short, the output from the rule system is a list of two query trees that correspond to these statements:

```
INSERT INTO shoelace_log VALUES (
    shoelace_data.sl_name, 6,
    current_user, current_timestamp )
FROM shoelace_data
WHERE 6 <> shoelace_data.sl_avail
    AND shoelace_data.sl_name = 'sl7';

UPDATE shoelace_data SET sl_avail = 6
WHERE sl_name = 'sl7';
```

These are executed in this order, and that is exactly what the rule was meant to do.

The substitutions and the added qualifications ensure that, if the original query would be, say:

```
UPDATE shoelace_data SET sl_color = 'green'
WHERE sl_name = 'sl7';
```

no log entry would get written. In that case, the original query tree does not contain a target list entry for sl_avail, so NEW.sl_avail will get replaced by shoelace_data.sl_avail. Thus, the extra command generated by the rule is:

```
INSERT INTO shoelace_log VALUES (  
    shoelace_data.sl_name, shoelace_data.sl_avail,
    current_user, current_timestamp )
FROM shoelace_data
WHERE shoelace_data.sl_avail <> shoelace_data.sl_avail
    AND shoelace_data.sl_name = 'sl7';
```

and that qualification will never be true.

It will also work if the original query modifies multiple rows. So if someone issued the command:
UPDATE shoelace_data SET sl_avail = 0
WHERE sl_color = 'black';

four rows in fact get updated (sl1, sl2, sl3, and sl4). But sl3 already has sl_avail = 0. In this case, the original query trees qualification is different and that results in the extra query tree:

INSERT INTO shoelace_log
SELECT shoelace_data.sl_name, 0,
    current_user, current_timestamp
FROM shoelace_data
WHERE 0 <> shoelace_data.sl_avail
    AND shoelace_data.sl_color = 'black';

being generated by the rule. This query tree will surely insert three new log entries. And that's absolutely correct.

Here we can see why it is important that the original query tree is executed last. If the UPDATE had been executed first, all the rows would have already been set to zero, so the logging INSERT would not find any row where 0 <> shoelace_data.sl_avail.

41.4.2. Cooperation with Views

A simple way to protect view relations from the mentioned possibility that someone can try to run INSERT, UPDATE, or DELETE on them is to let those query trees get thrown away. So we could create the rules:

CREATE RULE shoe_ins_protect AS ON INSERT TO shoe
    DO INSTEAD NOTHING;
CREATE RULE shoe_upd_protect AS ON UPDATE TO shoe
    DO INSTEAD NOTHING;
CREATE RULE shoe_del_protect AS ON DELETE TO shoe
    DO INSTEAD NOTHING;

If someone now tries to do any of these operations on the view relation shoe, the rule system will apply these rules. Since the rules have no actions and are INSTEAD, the resulting list of query trees will be empty and the whole query will become nothing because there is nothing left to be optimized or executed after the rule system is done with it.

A more sophisticated way to use the rule system is to create rules that rewrite the query tree into one that does the right operation on the real tables. To do that on the shoelace view, we create the following rules:

CREATE RULE shoelace_ins AS ON INSERT TO shoelace
    DO INSTEAD
    INSERT INTO shoelace_data VALUES (NEW.sl_name,
                                        NEW.sl_avail,
                                        NEW.sl_color,
                                        NEW.sl_len,
                                        NEW.sl_unit );

CREATE RULE shoelace_upd AS ON UPDATE TO shoelace
    DO INSTEAD
    UPDATE shoelace_data
        SET sl_name = NEW.sl_name,
The Rule System

```
sl_avail = NEW.sl_avail,
sl_color = NEW.sl_color,
sl_len = NEW.sl_len,
sl_unit = NEW.sl_unit
WHERE sl_name = OLD.sl_name;

CREATE RULE shoelace_del AS ON DELETE TO shoelace
  DO INSTEAD
    DELETE FROM shoelace_data
    WHERE sl_name = OLD.sl_name;

If you want to support RETURNING queries on the view, you need to make the rules include RE-
TURNING clauses that compute the view rows. This is usually pretty trivial for views on a single table,
but it's a bit tedious for join views such as shoelace. An example for the insert case is:

CREATE RULE shoelace_ins AS ON INSERT TO shoelace
  DO INSTEAD
    INSERT INTO shoelace_data VALUES (NEW.sl_name,
                                      NEW.sl_avail,
                                      NEW.sl_color,
                                      NEW.sl_len,
                                      NEW.sl_unit)
  RETURNING
    shoelace_data.*,
    (SELECT shoelace_data.sl_len * u.un_fact
     FROM unit u WHERE shoelace_data.sl_unit = u.un_name);

Note that this one rule supports both INSERT and INSERT RETURNING queries on the view — the
RETURNING clause is simply ignored for INSERT.

Now assume that once in a while, a pack of shoelaces arrives at the shop and a big parts list along
with it. But you don't want to manually update the shoelace view every time. Instead we set up
two little tables: one where you can insert the items from the part list, and one with a special trick.
The creation commands for these are:

CREATE TABLE shoelace_arrive (arr_name text,
                               arr_quant integer);

CREATE TABLE shoelace_ok (ok_name text,
                           ok_quant integer);

CREATE RULE shoelace_ok_ins AS ON INSERT TO shoelace_ok
  DO INSTEAD
    UPDATE shoelace
    SET sl_avail = sl_avail + NEW.ok_quant
    WHERE sl_name = NEW.ok_name;

Now you can fill the table shoelace_arrive with the data from the parts list:

SELECT * FROM shoelace_arrive;
```
Take a quick look at the current data:

```sql
SELECT * FROM shoelace;
```

```
<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl3</td>
<td>6</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl5</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl6</td>
<td>0</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>
```

(8 rows)

Now move the arrived shoelaces in:

```sql
INSERT INTO shoelace_ok SELECT * FROM shoelace_arrive;
```

and check the results:

```sql
SELECT * FROM shoelace ORDER BY sl_name;
```

```
<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl3</td>
<td>6</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl5</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl6</td>
<td>20</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>
```

(8 rows)

```sql
SELECT * FROM shoelace_log;
```

```
<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>log_who</th>
<th>log_when</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl7</td>
<td>6</td>
<td>Al</td>
<td>Tue Oct 20 19:14:45 1998 MET DST</td>
</tr>
<tr>
<td>sl3</td>
<td>10</td>
<td>Al</td>
<td>Tue Oct 20 19:25:16 1998 MET DST</td>
</tr>
<tr>
<td>sl6</td>
<td>20</td>
<td>Al</td>
<td>Tue Oct 20 19:25:16 1998 MET DST</td>
</tr>
<tr>
<td>sl8</td>
<td>21</td>
<td>Al</td>
<td>Tue Oct 20 19:25:16 1998 MET DST</td>
</tr>
</tbody>
</table>
```

(4 rows)

It's a long way from the one INSERT ... SELECT to these results. And the description of the
query-tree transformation will be the last in this chapter. First, there is the parser's output:
The Rule System

INSERT INTO shoelace_ok
SELECT shoelace_arrive.arr_name, shoelace_arrive.arr_quant
    FROM shoelace_arrive shoelace_arrive, shoelace_ok shoelace_ok;

Now the first rule shoelace_ok_ins is applied and turns this into:

UPDATE shoelace
    SET sl_avail = shoelace.sl_avail + shoelace_arrive.arr_quant
    FROM shoelace_arrive shoelace_arrive, shoelace_ok shoelace_ok,
        shoelace_ok old, shoelace_ok new,
        shoelace shoelace
    WHERE shoelace.sl_name = shoelace_arrive.arr_name;

and throws away the original INSERT on shoelace_ok. This rewritten query is passed to the rule system again, and the second applied rule shoelace_upd produces:

UPDATE shoelace_data
    SET sl_name = shoelace.sl_name,
        sl_avail = shoelace.sl_avail + shoelace_arrive.arr_quant,
        sl_color = shoelace.sl_color,
        sl_len = shoelace.sl_len,
        sl_unit = shoelace.sl_unit
    FROM shoelace_arrive shoelace_arrive, shoelace_ok shoelace_ok,
        shoelace_ok old, shoelace_ok new,
        shoelace shoelace, shoelace old,
        shoelace new, shoelace_data shoelace_data
    WHERE shoelace.sl_name = shoelace_arrive.arr_name
        AND shoelace_data.sl_name = shoelace.sl_name;

Again it's an INSTEAD rule and the previous query tree is trashed. Note that this query still uses the view shoelace. But the rule system isn't finished with this step, so it continues and applies the _RETURN rule on it, and we get:

UPDATE shoelace_data
    SET sl_name = s.sl_name,
        sl_avail = s.sl_avail + shoelace_arrive.arr_quant,
        sl_color = s.sl_color,
        sl_len = s.sl_len,
        sl_unit = s.sl_unit
    FROM shoelace_arrive shoelace_arrive, shoelace_ok shoelace_ok,
        shoelace_ok old, shoelace_ok new,
        shoelace shoelace, shoelace old,
        shoelace new, shoelace_data shoelace_data
        shoelace_data s, unit u
    WHERE s.sl_name = shoelace_arrive.arr_name
        AND shoelace_data.sl_name = s.sl_name;

Finally, the rule log_shoelace gets applied, producing the extra query tree:

INSERT INTO shoelace_log
SELECT s.sl_name,
    s.sl_avail + shoelace_arrive.arr_quant,
    current_user,
    current_timestamp

1148
The Rule System

FROM shoelace_arrive shoelace_arrive, shoelace_ok shoelace_ok,
  shoelace_ok old, shoelace_ok new,
  shoelace shoelace, shoelace old,
  shoelace new, shoelace_data shoelace_data,
  shoelace old, shoelace new,
  shoelace_data s, unit u,
  shoelace_data old, shoelace_data new
shoe lace_log shoe lace_log
WHERE s.sl_name = shoelace_arrive.arr_name
AND shoelace_data.sl_name = s.sl_name
AND (s.sl_avail + shoelace_arrive.arr_quant) <> s.sl_avail;

After that the rule system runs out of rules and returns the generated query trees.

So we end up with two final query trees that are equivalent to the SQL statements:

INSERT INTO shoelace_log
SELECT s.sl_name,
  s.sl_avail + shoelace_arrive.arr_quant,
  current_user,
  current_timestamp
FROM shoelace_arrive shoelace_arrive, shoelace_data
  shoelace_data,
  shoelace_data s
WHERE s.sl_name = shoelace_arrive.arr_name
AND shoelace_data.sl_name = s.sl_name
AND s.sl_avail + shoelace_arrive.arr_quant <> s.sl_avail;

UPDATE shoelace_data
SET sl_avail = shoelace_data.sl_avail +
  shoelace_arrive.arr_quant
FROM shoelace_arrive shoelace_arrive,
  shoelace_data shoelace_data,
  shoelace_data s
WHERE s.sl_name = shoelace_arrive.sl_name
AND shoelace_data.sl_name = s.sl_name;

The result is that data coming from one relation inserted into another, changed into updates on a third,
changed into updating a fourth plus logging that final update in a fifth gets reduced into two queries.

There is a little detail that's a bit ugly. Looking at the two queries, it turns out that the shoelace_data
relation appears twice in the range table where it could definitely be reduced to one. The planner
does not handle it and so the execution plan for the rule systems output of the INSERT will be

Nested Loop
  -> Merge Join
    -> Seq Scan
      -> Sort
      -> Seq Scan on s
    -> Seq Scan
      -> Sort
      -> Seq Scan on shoelace_arrive
  -> Seq Scan on shoelace_data

while omitting the extra range table entry would result in a

Merge Join
The Rule System

```
-> Seq Scan
   -> Sort
   -> Seq Scan on s
-> Seq Scan
   -> Sort
   -> Seq Scan on shoelace_arrive
```

which produces exactly the same entries in the log table. Thus, the rule system caused one extra scan on the table shoelace_data that is absolutely not necessary. And the same redundant scan is done once more in the UPDATE. But it was a really hard job to make that all possible at all.

Now we make a final demonstration of the PostgreSQL rule system and its power. Say you add some shoelaces with extraordinary colors to your database:

```
INSERT INTO shoelace VALUES ('sl9', 0, 'pink', 35.0, 'inch', 0.0);
INSERT INTO shoelace VALUES ('sl10', 1000, 'magenta', 40.0, 'inch', 0.0);
```

We would like to make a view to check which shoelace entries do not fit any shoe in color. The view for this is:

```
CREATE VIEW shoelace_mismatch AS
   SELECT * FROM shoelace WHERE NOT EXISTS
   (SELECT shoename FROM shoe WHERE slcolor = sl_color);
```

Its output is:

```
SELECT * FROM shoelace_mismatch;

sl_name | sl_avail | sl_color | sl_len | sl_unit | sl_len_cm
---------+----------+----------+--------+---------+-----------
sl9     |        0 | pink     |     35 | inch    |      88.9
sl10    |     1000 | magenta  |     40 | inch    |     101.6
```

Now we want to set it up so that mismatching shoelaces that are not in stock are deleted from the database. To make it a little harder for PostgreSQL, we don't delete it directly. Instead we create one more view:

```
CREATE VIEW shoelace_can_delete AS
   SELECT * FROM shoelace_mismatch WHERE sl_avail = 0;
```

and do it this way:

```
DELETE FROM shoelace WHERE EXISTS
   (SELECT * FROM shoelace_can_delete
    WHERE sl_name = shoelace.sl_name);
```

Voilà:

```
SELECT * FROM shoelace;

sl_name | sl_avail | sl_color | sl_len | sl_unit | sl_len_cm
---------+----------+----------+--------+---------+-----------
sl1     |        5 | black    |     80 | cm      |        80
```

1150
A DELETE on a view, with a subquery qualification that in total uses 4 nesting/joined views, where one of them itself has a subquery qualification containing a view and where calculated view columns are used, gets rewritten into one single query tree that deletes the requested data from a real table.

There are probably only a few situations out in the real world where such a construct is necessary. But it makes you feel comfortable that it works.

### 41.5. Rules and Privileges

Due to rewriting of queries by the PostgreSQL rule system, other tables/views than those used in the original query get accessed. When update rules are used, this can include write access to tables.

Rewrite rules don’t have a separate owner. The owner of a relation (table or view) is automatically the owner of the rewrite rules that are defined for it. The PostgreSQL rule system changes the behavior of the default access control system. Relations that are used due to rules get checked against the privileges of the rule owner, not the user invoking the rule. This means that users only need the required privileges for the tables/views that are explicitly named in their queries.

For example: A user has a list of phone numbers where some of them are private, the others are of interest for the assistant of the office. The user can construct the following:

```
CREATE TABLE phone_data (person text, phone text, private boolean);
CREATE VIEW phone_number AS
    SELECT person, CASE WHEN NOT private THEN phone END AS phone
    FROM phone_data;
GRANT SELECT ON phone_number TO assistant;
```

Nobody except that user (and the database superusers) can access the `phone_data` table. But because of the GRANT, the assistant can run a SELECT on the `phone_number` view. The rule system will rewrite the SELECT from `phone_number` into a SELECT from `phone_data`. Since the user is the owner of `phone_number` and therefore the owner of the rule, the read access to `phone_data` is now checked against the user's privileges and the query is permitted. The check for accessing `phone_number` is also performed, but this is done against the invoking user, so nobody but the user and the assistant can use it.

The privileges are checked rule by rule. So the assistant is for now the only one who can see the public phone numbers. But the assistant can set up another view and grant access to that to the public. Then, anyone can see the `phone_number` data through the assistant's view. What the assistant cannot do is to create a view that directly accesses `phone_data`. (Actually the assistant can, but it will not work since every access will be denied during the permission checks.) And as soon as the user notices that the assistant opened their `phone_number` view, the user can revoke the assistant's access. Immediately, any access to the assistant's view would fail.

One might think that this rule-by-rule checking is a security hole, but in fact it isn't. But if it did not work this way, the assistant could set up a table with the same columns as `phone_number` and copy the data to there once per day. Then it’s the assistant's own data and the assistant can grant access to everyone they want. A GRANT command means, “I trust you”. If someone you trust does the thing above, it's time to think it over and then use REVOKE.
Note that while views can be used to hide the contents of certain columns using the technique shown above, they cannot be used to reliably conceal the data in unseen rows unless the `security_barrier` flag has been set. For example, the following view is insecure:

```
CREATE VIEW phone_number AS
    SELECT person, phone FROM phone_data WHERE phone NOT LIKE '412%';
```

This view might seem secure, since the rule system will rewrite any `SELECT` from `phone_number` into a `SELECT` from `phone_data` and add the qualification that only entries where phone does not begin with 412 are wanted. But if the user can create their own functions, it is not difficult to convince the planner to execute the user-defined function prior to the `NOT LIKE` expression. For example:

```
CREATE FUNCTION tricky(text, text) RETURNS bool AS $$
BEGIN
    RAISE NOTICE '% => %', $1, $2;
    RETURN true;
END$$ LANGUAGE plpgsql COST 0.0000000000000000000001;
SELECT * FROM phone_number WHERE tricky(person, phone);
```

Every person and phone number in the `phone_data` table will be printed as a NOTICE, because the planner will choose to execute the inexpensive `tricky` function before the more expensive `NOT LIKE`. Even if the user is prevented from defining new functions, built-in functions can be used in similar attacks. (For example, most casting functions include their input values in the error messages they produce.)

Similar considerations apply to update rules. In the examples of the previous section, the owner of the tables in the example database could grant the privileges `SELECT`, `INSERT`, `UPDATE`, and `DELETE` on the `shoelace` view to someone else, but only `SELECT` on `shoelace_log`. The rule action to write log entries will still be executed successfully, and that other user could see the log entries. But they could not create fake entries, nor could they manipulate or remove existing ones. In this case, there is no possibility of subverting the rules by convincing the planner to alter the order of operations, because the only rule which references `shoelace_log` is an unqualified `INSERT`. This might not be true in more complex scenarios.

When it is necessary for a view to provide row level security, the `security_barrier` attribute should be applied to the view. This prevents maliciously-chosen functions and operators from being passed values from rows until after the view has done its work. For example, if the view shown above had been created like this, it would be secure:

```
CREATE VIEW phone_number WITH (security_barrier) AS
    SELECT person, phone FROM phone_data WHERE phone NOT LIKE '412%';
```

Views created with the `security_barrier` may perform far worse than views created without this option. In general, there is no way to avoid this: the fastest possible plan must be rejected if it may compromise security. For this reason, this option is not enabled by default.

The query planner has more flexibility when dealing with functions that have no side effects. Such functions are referred to as `LEAKPROOF`, and include many simple, commonly used operators, such as many equality operators. The query planner can safely allow such functions to be evaluated at any point in the query execution process, since invoking them on rows invisible to the user will not leak any information about the unseen rows. Further, functions which do not take arguments or which are not passed any arguments from the security barrier view do not have to be marked as `LEAKPROOF` to
be pushed down, as they never receive data from the view. In contrast, a function that might throw an error depending on the values received as arguments (such as one that throws an error in the event of overflow or division by zero) is not leak-proof, and could provide significant information about the unseen rows if applied before the security view’s row filters.

It is important to understand that even a view created with the security_barrier option is intended to be secure only in the limited sense that the contents of the invisible tuples will not be passed to possibly-insecure functions. The user may well have other means of making inferences about the unseen data; for example, they can see the query plan using EXPLAIN, or measure the run time of queries against the view. A malicious attacker might be able to infer something about the amount of unseen data, or even gain some information about the data distribution or most common values (since these things may affect the run time of the plan; or even, since they are also reflected in the optimizer statistics, the choice of plan). If these types of “covert channel” attacks are of concern, it is probably unwise to grant any access to the data at all.

41.6. Rules and Command Status

The PostgreSQL server returns a command status string, such as INSERT 149592 1, for each command it receives. This is simple enough when there are no rules involved, but what happens when the query is rewritten by rules?

Rules affect the command status as follows:

• If there is no unconditional INSTEAD rule for the query, then the originally given query will be executed, and its command status will be returned as usual. (But note that if there were any conditional INSTEAD rules, the negation of their qualifications will have been added to the original query. This might reduce the number of rows it processes, and if so the reported status will be affected.)

• If there is any unconditional INSTEAD rule for the query, then the original query will not be executed at all. In this case, the server will return the command status for the last query that was inserted by an INSTEAD rule (conditional or unconditional) and is of the same command type (INSERT, UPDATE, or DELETE) as the original query. If no query meeting those requirements is added by any rule, then the returned command status shows the original query type and zeroes for the row-count and OID fields.

The programmer can ensure that any desired INSTEAD rule is the one that sets the command status in the second case, by giving it the alphabetically last rule name among the active rules, so that it gets applied last.

41.7. Rules Versus Triggers

Many things that can be done using triggers can also be implemented using the PostgreSQL rule system. One of the things that cannot be implemented by rules are some kinds of constraints, especially foreign keys. It is possible to place a qualified rule that rewrites a command to NOTHING if the value of a column does not appear in another table. But then the data is silently thrown away and that's not a good idea. If checks for valid values are required, and in the case of an invalid value an error message should be generated, it must be done by a trigger.

In this chapter, we focused on using rules to update views. All of the update rule examples in this chapter can also be implemented using INSTEAD OF triggers on the views. Writing such triggers is often easier than writing rules, particularly if complex logic is required to perform the update.

For the things that can be implemented by both, which is best depends on the usage of the database. A trigger is fired once for each affected row. A rule modifies the query or generates an additional query. So if many rows are affected in one statement, a rule issuing one extra command is likely to be faster than a trigger that is called for every single row and must re-determine what to do many times. However, the trigger approach is conceptually far simpler than the rule approach, and is easier for novices to get right.
Here we show an example of how the choice of rules versus triggers plays out in one situation. There are two tables:

```sql
CREATE TABLE computer (
    hostname        text,    -- indexed
    manufacturer    text     -- indexed
);

CREATE TABLE software (
    software        text,    -- indexed
    hostname        text     -- indexed
);
```

Both tables have many thousands of rows and the indexes on `hostname` are unique. The rule or trigger should implement a constraint that deletes rows from `software` that reference a deleted computer. The trigger would use this command:

```sql
DELETE FROM software WHERE hostname = $1;
```

Since the trigger is called for each individual row deleted from `computer`, it can prepare and save the plan for this command and pass the `hostname` value in the parameter. The rule would be written as:

```sql
CREATE RULE computer_del AS ON DELETE TO computer
    DO DELETE FROM software WHERE hostname = OLD.hostname;
```

Now we look at different types of deletes. In the case of a:

```sql
DELETE FROM computer WHERE hostname = 'mypc.local.net';
```

the table `computer` is scanned by index (fast), and the command issued by the trigger would also use an index scan (also fast). The extra command from the rule would be:

```sql
DELETE FROM software WHERE computer.hostname = 'mypc.local.net'
    AND software.hostname = computer.hostname;
```

Since there are appropriate indexes set up, the planner will create a plan of

```
Nestloop
  -> Index Scan using comp_hostidx on computer
  -> Index Scan using soft_hostidx on software
```

So there would be not that much difference in speed between the trigger and the rule implementation.

With the next delete we want to get rid of all the 2000 computers where the `hostname` starts with `old`. There are two possible commands to do that. One is:

```sql
DELETE FROM computer WHERE hostname >= 'old'
    AND hostname <  'ole'
```

The command added by the rule will be:

```sql
DELETE FROM software WHERE computer.hostname >= 'old' AND
    computer.hostname <  'ole'
```
The Rule System

AND software.hostname = computer.hostname;

with the plan

Hash Join
  -> Seq Scan on software
  -> Hash
  -> Index Scan using comp_hostidx on computer

The other possible command is:

DELETE FROM computer WHERE hostname ~ '^old';

which results in the following executing plan for the command added by the rule:

Nestloop
  -> Index Scan using comp_hostidx on computer
  -> Index Scan using soft_hostidx on software

This shows, that the planner does not realize that the qualification for hostname in computer could also be used for an index scan on software when there are multiple qualification expressions combined with AND, which is what it does in the regular-expression version of the command. The trigger will get invoked once for each of the 2000 old computers that have to be deleted, and that will result in one index scan over computer and 2000 index scans over software. The rule implementation will do it with two commands that use indexes. And it depends on the overall size of the table software whether the rule will still be faster in the sequential scan situation. 2000 command executions from the trigger over the SPI manager take some time, even if all the index blocks will soon be in the cache.

The last command we look at is:

DELETE FROM computer WHERE manufacturer = 'bim';

Again this could result in many rows to be deleted from computer. So the trigger will again run many commands through the executor. The command generated by the rule will be:

DELETE FROM software WHERE computer.manufacturer = 'bim'
  AND software.hostname = computer.hostname;

The plan for that command will again be the nested loop over two index scans, only using a different index on computer:

Nestloop
  -> Index Scan using comp_manufidx on computer
  -> Index Scan using soft_hostidx on software

In any of these cases, the extra commands from the rule system will be more or less independent from the number of affected rows in a command.

The summary is, rules will only be significantly slower than triggers if their actions result in large and badly qualified joins, a situation where the planner fails.
PostgreSQL allows user-defined functions to be written in other languages besides SQL and C. These other languages are generically called *procedural languages* (PLs). For a function written in a procedural language, the database server has no built-in knowledge about how to interpret the function's source text. Instead, the task is passed to a special handler that knows the details of the language. The handler could either do all the work of parsing, syntax analysis, execution, etc. itself, or it could serve as “glue” between PostgreSQL and an existing implementation of a programming language. The handler itself is a C language function compiled into a shared object and loaded on demand, just like any other C function.

There are currently four procedural languages available in the standard PostgreSQL distribution: PL/pgSQL (Chapter 43), PL/Tcl (Chapter 44), PL/Perl (Chapter 45), and PL/Python (Chapter 46). There are additional procedural languages available that are not included in the core distribution. Appendix H has information about finding them. In addition other languages can be defined by users; the basics of developing a new procedural language are covered in Chapter 56.

### 42.1. Installing Procedural Languages

A procedural language must be “installed” into each database where it is to be used. But procedural languages installed in the database `template1` are automatically available in all subsequently created databases, since their entries in `template1` will be copied by `CREATE DATABASE`. So the database administrator can decide which languages are available in which databases and can make some languages available by default if desired.

For the languages supplied with the standard distribution, it is only necessary to execute `CREATE EXTENSION language_name` to install the language into the current database. The manual procedure described below is only recommended for installing languages that have not been packaged as extensions.

#### Manual Procedural Language Installation

A procedural language is installed in a database in five steps, which must be carried out by a database superuser. In most cases the required SQL commands should be packaged as the installation script of an “extension”, so that `CREATE EXTENSION` can be used to execute them.

1. The shared object for the language handler must be compiled and installed into an appropriate library directory. This works in the same way as building and installing modules with regular user-defined C functions does; see Section 38.10.5. Often, the language handler will depend on an external library that provides the actual programming language engine; if so, that must be installed as well.

2. The handler must be declared with the command

   ```sql
   CREATE FUNCTION handler_function_name()
   RETURNS language_handler
   AS 'path-to-shared-object'
   LANGUAGE C;
   ```

   The special return type of `language_handler` tells the database system that this function does not return one of the defined SQL data types and is not directly usable in SQL statements.

3. (Optional) Optionally, the language handler can provide an “inline” handler function that executes anonymous code blocks (DO commands) written in this language. If an inline handler function is provided by the language, declare it with a command like
CREATE FUNCTION inline_function_name(internal)
    RETURNS void
    AS 'path-to-shared-object'
    LANGUAGE C;

4. (Optional) Optionally, the language handler can provide a “validator” function that checks a function definition for correctness without actually executing it. The validator function is called by CREATE FUNCTION if it exists. If a validator function is provided by the language, declare it with a command like

CREATE FUNCTION validator_function_name(oid)
    RETURNS void
    AS 'path-to-shared-object'
    LANGUAGE C STRICT;

5. Finally, the PL must be declared with the command

CREATE [TRUSTED] [PROCEDURAL] LANGUAGE language-name
    HANDLER handler_function_name
    [INLINE inline_function_name]
    [VALIDATOR validator_function_name] ;

The optional key word TRUSTED specifies that the language does not grant access to data that the user would not otherwise have. Trusted languages are designed for ordinary database users (those without superuser privilege) and allows them to safely create functions and procedures. Since PL functions are executed inside the database server, the TRUSTED flag should only be given for languages that do not allow access to database server internals or the file system. The languages PL/pgSQL, PL/Tcl, and PL/Perl are considered trusted; the languages PL/TclU, PL/PerlU, and PL/PythonU are designed to provide unlimited functionality and should not be marked trusted.

Example 42.1 shows how the manual installation procedure would work with the language PL/Perl.

**Example 42.1. Manual Installation of PL/Perl**

The following command tells the database server where to find the shared object for the PL/Perl language's call handler function:

CREATE FUNCTION plperl_call_handler() RETURNS language_handler AS
    '$libdir/plperl' LANGUAGE C;

PL/Perl has an inline handler function and a validator function, so we declare those too:

CREATE FUNCTION plperl_inline_handler(internal) RETURNS void AS
    '$libdir/plperl' LANGUAGE C;

CREATE FUNCTION plperl_validator(oid) RETURNS void AS
    '$libdir/plperl' LANGUAGE C STRICT;

The command:

CREATE TRUSTED PROCEDURAL LANGUAGE plperl
    HANDLER plperl_call_handler
    INLINE plperl_inline_handler
    VALIDATOR plperl_validator;
then defines that the previously declared functions should be invoked for functions and procedures where the language attribute is \texttt{plperl}.

In a default PostgreSQL installation, the handler for the PL/pgSQL language is built and installed into the “library” directory; furthermore, the PL/pgSQL language itself is installed in all databases. If Tcl support is configured in, the handlers for PL/Tcl and PL/TclU are built and installed in the library directory, but the language itself is not installed in any database by default. Likewise, the PL/Perl and PL/PerlU handlers are built and installed if Perl support is configured, and the PL/PythonU handler is installed if Python support is configured, but these languages are not installed by default.
Chapter 43. PL/pgSQL - SQL Procedural Language

43.1. Overview

PL/pgSQL is a loadable procedural language for the PostgreSQL database system. The design goals of PL/pgSQL were to create a loadable procedural language that

• can be used to create functions and triggers,
• adds control structures to the SQL language,
• can perform complex computations,
• inherits all user-defined types, functions, and operators,
• can be defined to be trusted by the server,
• is easy to use.

Functions created with PL/pgSQL can be used anywhere that built-in functions could be used. For example, it is possible to create complex conditional computation functions and later use them to define operators or use them in index expressions.

In PostgreSQL 9.0 and later, PL/pgSQL is installed by default. However it is still a loadable module, so especially security-conscious administrators could choose to remove it.

43.1.1. Advantages of Using PL/pgSQL

SQL is the language PostgreSQL and most other relational databases use as query language. It's portable and easy to learn. But every SQL statement must be executed individually by the database server.

That means that your client application must send each query to the database server, wait for it to be processed, receive and process the results, do some computation, then send further queries to the server. All this incurs interprocess communication and will also incur network overhead if your client is on a different machine than the database server.

With PL/pgSQL you can group a block of computation and a series of queries inside the database server, thus having the power of a procedural language and the ease of use of SQL, but with considerable savings of client/server communication overhead.

• Extra round trips between client and server are eliminated
• Intermediate results that the client does not need do not have to be marshaled or transferred between server and client
• Multiple rounds of query parsing can be avoided

This can result in a considerable performance increase as compared to an application that does not use stored functions.

Also, with PL/pgSQL you can use all the data types, operators and functions of SQL.

43.1.2. Supported Argument and Result Data Types

Functions written in PL/pgSQL can accept as arguments any scalar or array data type supported by the server, and they can return a result of any of these types. They can also accept or return any composite
type (row type) specified by name. It is also possible to declare a PL/pgSQL function as accepting record, which means that any composite type will do as input, or as returning record, which means that the result is a row type whose columns are determined by specification in the calling query, as discussed in Section 7.2.1.4.

PL/pgSQL functions can be declared to accept a variable number of arguments by using the VARIADIC marker. This works exactly the same way as for SQL functions, as discussed in Section 38.5.5.

PL/pgSQL functions can also be declared to accept and return the polymorphic types anyelement, anyarray, anynonarray, anyenum, and anyrange. The actual data types handled by a polymorphic function can vary from call to call, as discussed in Section 38.2.5. An example is shown in Section 43.3.1.

PL/pgSQL functions can also be declared to return a “set” (or table) of any data type that can be returned as a single instance. Such a function generates its output by executing RETURN NEXT for each desired element of the result set, or by using RETURN QUERY to output the result of evaluating a query.

Finally, a PL/pgSQL function can be declared to return void if it has no useful return value. (Alternatively, it could be written as a procedure in that case.)

PL/pgSQL functions can also be declared with output parameters in place of an explicit specification of the return type. This does not add any fundamental capability to the language, but it is often convenient, especially for returning multiple values. The RETURNS TABLE notation can also be used in place of RETURNS SETOF.

Specific examples appear in Section 43.3.1 and Section 43.6.1.

43.2. Structure of PL/pgSQL

Functions written in PL/pgSQL are defined to the server by executing CREATE FUNCTION commands. Such a command would normally look like, say,

```
CREATE FUNCTION somefunc(integer, text) RETURNS integer
AS 'function body text'
LANGUAGE plpgsql;
```

The function body is simply a string literal so far as CREATE FUNCTION is concerned. It is often helpful to use dollar quoting (see Section 4.1.2.4) to write the function body, rather than the normal single quote syntax. Without dollar quoting, any single quotes or backslashes in the function body must be escaped by doubling them. Almost all the examples in this chapter use dollar-quoted literals for their function bodies.

PL/pgSQL is a block-structured language. The complete text of a function body must be a block. A block is defined as:

```
[ <<label>> ]
[ DECLARE
declarations ]
BEGIN
  statements
END [ label ];
```

Each declaration and each statement within a block is terminated by a semicolon. A block that appears within another block must have a semicolon after END, as shown above; however the final END that concludes a function body does not require a semicolon.
A common mistake is to write a semicolon immediately after BEGIN. This is incorrect and will result in a syntax error.

A **label** is only needed if you want to identify the block for use in an EXIT statement, or to qualify the names of the variables declared in the block. If a label is given after END, it must match the label at the block's beginning.

All key words are case-insensitive. Identifiers are implicitly converted to lower case unless double-quoted, just as they are in ordinary SQL commands.

Comments work the same way in PL/pgSQL code as in ordinary SQL. A double dash (--) starts a comment that extends to the end of the line. A /* starts a block comment that extends to the matching occurrence of */. Block comments nest.

Any statement in the statement section of a block can be a **subblock**. Subblocks can be used for logical grouping or to localize variables to a small group of statements. Variables declared in a subblock mask any similarly-named variables of outer blocks for the duration of the subblock; but you can access the outer variables anyway if you qualify their names with their block's label. For example:

```sql
CREATE FUNCTION somefunc() RETURNS integer AS $$
<< outerblock >>
DECLARE
    quantity integer := 30;
BEGIN
    RAISE NOTICE 'Quantity here is %', quantity;  -- Prints 30
    quantity := 50;
    --
    -- Create a subblock
    --
    DECLARE
        quantity integer := 80;
    BEGIN
        RAISE NOTICE 'Quantity here is %', quantity;  -- Prints 80
        RAISE NOTICE 'Outer quantity here is %', outerblock.quantity;  -- Prints 50
    END;

    RAISE NOTICE 'Quantity here is %', quantity;  -- Prints 50

    RETURN quantity;
END;
$$ LANGUAGE plpgsql;
```

There is actually a hidden “outer block” surrounding the body of any PL/pgSQL function. This block provides the declarations of the function's parameters (if any), as well as some special variables such as FOUND (see Section 43.5.5). The outer block is labeled with the function's name, meaning that parameters and special variables can be qualified with the function's name.

It is important not to confuse the use of BEGIN/END for grouping statements in PL/pgSQL with the similarly-named SQL commands for transaction control. PL/pgSQL's BEGIN/END are only for group-
PL/pgSQL - SQL Pro-
cedural Language

ing; they do not start or end a transaction. See Section 43.8 for information on managing transactions in PL/pgSQL. Also, a block containing an EXCEPTION clause effectively forms a subtransaction that can be rolled back without affecting the outer transaction. For more about that see Section 43.6.8.

43.3. Declarations

All variables used in a block must be declared in the declarations section of the block. (The only exceptions are that the loop variable of a FOR loop iterating over a range of integer values is automatically declared as an integer variable, and likewise the loop variable of a FOR loop iterating over a cursor's result is automatically declared as a record variable.)

PL/pgSQL variables can have any SQL data type, such as integer, varchar, and char.

Here are some examples of variable declarations:

```
user_id integer;
quantity numeric(5);
url varchar;
myrow tablename%ROWTYPE;
myfield tablename.columnname%TYPE;
arow RECORD;
```

The general syntax of a variable declaration is:

```
name [ CONSTANT ] type [ COLLATE collation_name ] [ NOT NULL ]
[ { DEFAULT | := | = } expression ];
```

The DEFAULT clause, if given, specifies the initial value assigned to the variable when the block is entered. If the DEFAULT clause is not given then the variable is initialized to the SQL null value. The CONSTANT option prevents the variable from being assigned to after initialization, so that its value will remain constant for the duration of the block. The COLLATE option specifies a collation to use for the variable (see Section 43.3.6). If NOT NULL is specified, an assignment of a null value results in a run-time error. All variables declared as NOT NULL must have a nonnull default value specified.

Equal (=) can be used instead of PL/SQL-compliant :=.

A variable's default value is evaluated and assigned to the variable each time the block is entered (not just once per function call). So, for example, assigning now() to a variable of type timestamp causes the variable to have the time of the current function call, not the time when the function was precompiled.

Examples:

```
quantity integer DEFAULT 32;
url varchar := 'http://mysite.com';
user_id CONSTANT integer := 10;
```

43.3.1. Declaring Function Parameters

Parameters passed to functions are named with the identifiers $1, $2, etc. Optionally, aliases can be declared for $n parameter names for increased readability. Either the alias or the numeric identifier can then be used to refer to the parameter value.

There are two ways to create an alias. The preferred way is to give a name to the parameter in the CREATE FUNCTION command, for example:
CREATE FUNCTION sales_tax(subtotal real) RETURNS real AS $$
BEGIN
    RETURN subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;

The other way is to explicitly declare an alias, using the declaration syntax

ame ALIAS FOR $n;

The same example in this style looks like:

CREATE FUNCTION sales_tax(real) RETURNS real AS $$
DECLARE
    subtotal ALIAS FOR $1;
BEGIN
    RETURN subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;

Note

These two examples are not perfectly equivalent. In the first case, subtotal could be referenced as sales_tax.subtotal, but in the second case it could not. (Had we attached a label to the inner block, subtotal could be qualified with that label instead.)

Some more examples:

CREATE FUNCTION instr(varchar, integer) RETURNS integer AS $$
DECLARE
    v_string ALIAS FOR $1;
    index ALIAS FOR $2;
BEGIN
    -- some computations using v_string and index here
END;
$$ LANGUAGE plpgsql;

CREATE FUNCTION concat_selected_fields(in_t sometablename) RETURNS text AS $$
BEGIN
    RETURN in_t.f1 || in_t.f3 || in_t.f5 || in_t.f7;
END;
$$ LANGUAGE plpgsql;

When a PL/pgSQL function is declared with output parameters, the output parameters are given $n names and optional aliases in just the same way as the normal input parameters. An output parameter is effectively a variable that starts out NULL; it should be assigned to during the execution of the function. The final value of the parameter is what is returned. For instance, the sales-tax example could also be done this way:

CREATE FUNCTION sales_tax(subtotal real, OUT tax real) AS $$
BEGIN

1163
tax := subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;

Notice that we omitted RETURNS real — we could have included it, but it would be redundant.

Output parameters are most useful when returning multiple values. A trivial example is:

CREATE FUNCTION sum_n_product(x int, y int, OUT sum int, OUT prod int) AS $$
BEGIN
  sum := x + y;
  prod := x * y;
END;
$$ LANGUAGE plpgsql;

As discussed in Section 38.5.4, this effectively creates an anonymous record type for the function's results. If a RETURNS clause is given, it must say RETURNS record.

Another way to declare a PL/pgSQL function is with RETURNS TABLE, for example:

CREATE FUNCTION extended_sales(p_itemno int)
RETURNS TABLE(quantity int, total numeric) AS $$
BEGIN
  RETURN QUERY SELECT s.quantity, s.quantity * s.price FROM sales AS s
  WHERE s.itemno = p_itemno;
END;
$$ LANGUAGE plpgsql;

This is exactly equivalent to declaring one or more OUT parameters and specifying RETURNS SETOF sometype.

When the return type of a PL/pgSQL function is declared as a polymorphic type (anetype, anyarray, anynonarray, anyenum, or anyrange), a special parameter $0 is created. Its data type is the actual return type of the function, as deduced from the actual input types (see Section 38.2.5). This allows the function to access its actual return type as shown in Section 43.3.3. $0 is initialized to null and can be modified by the function, so it can be used to hold the return value if desired, though that is not required. $0 can also be given an alias. For example, this function works on any data type that has a + operator:

CREATE FUNCTION add_three_values(v1 anyelement, v2 anyelement, v3 anyelement)
RETURNS anyelement AS $$
DECLARE
  result ALIAS FOR $0;
BEGIN
  result := v1 + v2 + v3;
  RETURN result;
END;
$$ LANGUAGE plpgsql;

The same effect can be obtained by declaring one or more output parameters as polymorphic types. In this case the special $0 parameter is not used; the output parameters themselves serve the same purpose. For example:
CREATE FUNCTION add_three_values(v1 anyelement, v2 anyelement, v3 anyelement, OUT sum anyelement) AS $$
BEGIN
  sum := v1 + v2 + v3;
END;
$$ LANGUAGE plpgsql;

43.3.2. ALIAS

newname ALIAS FOR oldname;

The ALIAS syntax is more general than is suggested in the previous section; you can declare an alias for any variable, not just function parameters. The main practical use for this is to assign a different name for variables with predetermined names, such as NEW or OLD within a trigger function.

Examples:

DECLARE
  prior ALIAS FOR old;
  updated ALIAS FOR new;

Since ALIAS creates two different ways to name the same object, unrestricted use can be confusing. It's best to use it only for the purpose of overriding predetermined names.

43.3.3. Copying Types

variable%TYPE

%TYPE provides the data type of a variable or table column. You can use this to declare variables that will hold database values. For example, let's say you have a column named user_id in your users table. To declare a variable with the same data type as users.user_id you write:

user_id users.user_id%TYPE;

By using %TYPE you don't need to know the data type of the structure you are referencing, and most importantly, if the data type of the referenced item changes in the future (for instance: you change the type of user_id from integer to real), you might not need to change your function definition.

%TYPE is particularly valuable in polymorphic functions, since the data types needed for internal variables can change from one call to the next. Appropriate variables can be created by applying %TYPE to the function's arguments or result placeholders.

43.3.4. Row Types

name table_name%ROWTYPE;
name composite_type_name;

A variable of a composite type is called a row variable (or row-type variable). Such a variable can hold a whole row of a SELECT or FOR query result, so long as that query's column set matches the declared type of the variable. The individual fields of the row value are accessed using the usual dot notation, for example rowvar.field.
A row variable can be declared to have the same type as the rows of an existing table or view, by using the `table_name%ROWTYPE` notation; or it can be declared by giving a composite type's name. (Since every table has an associated composite type of the same name, it actually does not matter in PostgreSQL whether you write `%ROWTYPE` or not. But the form with `%ROWTYPE` is more portable.)

Parameters to a function can be composite types (complete table rows). In that case, the corresponding identifier $n$ will be a row variable, and fields can be selected from it, for example `$1.user_id`.

Here is an example of using composite types. `table1` and `table2` are existing tables having at least the mentioned fields:

```sql
CREATE FUNCTION merge_fields(t_row table1) RETURNS text AS $$
DECLARE
   t2_row table2%ROWTYPE;
BEGIN
   SELECT * INTO t2_row FROM table2 WHERE ... ;
   RETURN t_row.f1 || t2_row.f3 || t_row.f5 || t2_row.f7;
END;
$$ LANGUAGE plpgsql;

SELECT merge_fields(t.*) FROM table1 t WHERE ... ;
```

### 43.3.5. Record Types

Record variables are similar to row-type variables, but they have no predefined structure. They take on the actual row structure of the row they are assigned during a `SELECT` or `FOR` command. The substructure of a record variable can change each time it is assigned to. A consequence of this is that until a record variable is first assigned to, it has no substructure, and any attempt to access a field in it will draw a run-time error.

Note that `RECORD` is not a true data type, only a placeholder. One should also realize that when a PL/pgSQL function is declared to return type `record`, this is not quite the same concept as a record variable, even though such a function might use a record variable to hold its result. In both cases the actual row structure is unknown when the function is written, but for a function returning `record` the actual structure is determined when the calling query is parsed, whereas a record variable can change its row structure on-the-fly.

### 43.3.6. Collation of PL/pgSQL Variables

When a PL/pgSQL function has one or more parameters of collatable data types, a collation is identified for each function call depending on the collations assigned to the actual arguments, as described in Section 23.2. If a collation is successfully identified (i.e., there are no conflicts of implicit collations among the arguments) then all the collatable parameters are treated as having that collation implicitly. This will affect the behavior of collation-sensitive operations within the function. For example, consider

```sql
CREATE FUNCTION less_than(a text, b text) RETURNS boolean AS $$
BEGIN
   RETURN a < b;
END;
$$ LANGUAGE plpgsql;

SELECT less_than(text_field_1, text_field_2) FROM table1;
```
SELECT less_than(text_field_1, text_field_2 COLLATE "C") FROM table1;

The first use of less_than will use the common collation of text_field_1 and text_field_2 for the comparison, while the second use will use C collation.

Furthermore, the identified collation is also assumed as the collation of any local variables that are of collatable types. Thus this function would not work any differently if it were written as

CREATE FUNCTION less_than(a text, b text) RETURNS boolean AS $$
DECLARE
    local_a text := a;
    local_b text := b;
BEGIN
    RETURN local_a < local_b;
END;
$$ LANGUAGE plpgsql;

If there are no parameters of collatable data types, or no common collation can be identified for them, then parameters and local variables use the default collation of their data type (which is usually the database's default collation, but could be different for variables of domain types).

A local variable of a collatable data type can have a different collation associated with it by including the COLLATE option in its declaration, for example

DECLARE
    local_a text COLLATE "en_US";

This option overrides the collation that would otherwise be given to the variable according to the rules above.

Also, of course explicit COLLATE clauses can be written inside a function if it is desired to force a particular collation to be used in a particular operation. For example,

CREATE FUNCTION less_than_c(a text, b text) RETURNS boolean AS $$
BEGIN
    RETURN a < b COLLATE "C";
END;
$$ LANGUAGE plpgsql;

This overrides the collations associated with the table columns, parameters, or local variables used in the expression, just as would happen in a plain SQL command.

### 43.4. Expressions

All expressions used in PL/pgSQL statements are processed using the server's main SQL executor. For example, when you write a PL/pgSQL statement like

IF expression THEN ...

PL/pgSQL will evaluate the expression by feeding a query like

SELECT expression

to the main SQL engine. While forming the SELECT command, any occurrences of PL/pgSQL variable names are replaced by parameters, as discussed in detail in Section 43.11.1. This allows the query
plan for the SELECT to be prepared just once and then reused for subsequent evaluations with different values of the variables. Thus, what really happens on first use of an expression is essentially a PREPARE command. For example, if we have declared two integer variables \( x \) and \( y \), and we write

\[
\text{IF } x < y \text{ THEN } ...
\]

what happens behind the scenes is equivalent to

\[
\text{PREPARE statement_name(integer, integer) AS SELECT $1 < $2;}
\]

and then this prepared statement is EXECUTEd for each execution of the IF statement, with the current values of the PL/pgSQL variables supplied as parameter values. Normally these details are not important to a PL/pgSQL user, but they are useful to know when trying to diagnose a problem. More information appears in Section 43.11.2.

### 43.5. Basic Statements

In this section and the following ones, we describe all the statement types that are explicitly understood by PL/pgSQL. Anything not recognized as one of these statement types is presumed to be an SQL command and is sent to the main database engine to execute, as described in Section 43.5.2 and Section 43.5.3.

#### 43.5.1. Assignment

An assignment of a value to a PL/pgSQL variable is written as:

\[
\text{variable} \{ := | = \} \text{expression;}
\]

As explained previously, the expression in such a statement is evaluated by means of an SQL SELECT command sent to the main database engine. The expression must yield a single value (possibly a row value, if the variable is a row or record variable). The target variable can be a simple variable (optionally qualified with a block name), a field of a row or record variable, or an element of an array that is a simple variable or field. Equal (\( = \)) can be used instead of PL/SQL-compliant \( := \).

If the expression's result data type doesn't match the variable's data type, the value will be coerced as though by an assignment cast (see Section 10.4). If no assignment cast is known for the pair of data types involved, the PL/pgSQL interpreter will attempt to convert the result value textually, that is by applying the result type's output function followed by the variable type's input function. Note that this could result in run-time errors generated by the input function, if the string form of the result value is not acceptable to the input function.

Examples:

\[
\text{tax := subtotal * 0.06;}
\]
\[
\text{my_record.user_id := 20;}
\]

#### 43.5.2. Executing a Command With No Result

For any SQL command that does not return rows, for example INSERT without a RETURNING clause, you can execute the command within a PL/pgSQL function just by writing the command. Any PL/pgSQL variable name appearing in the command text is treated as a parameter, and then the current value of the variable is provided as the parameter value at runtime. This is exactly like the processing described earlier for expressions; for details see Section 43.11.1.
When executing a SQL command in this way, PL/pgSQL may cache and re-use the execution plan for the command, as discussed in Section 43.11.2.

Sometimes it is useful to evaluate an expression or `SELECT` query but discard the result, for example when calling a function that has side-effects but no useful result value. To do this in PL/pgSQL, use the `PERFORM` statement:

```
PERFORM query;
```

This executes `query` and discards the result. Write the `query` the same way you would write an SQL `SELECT` command, but replace the initial keyword `SELECT` with `PERFORM`. For `WITH` queries, use `PERFORM` and then place the query in parentheses. (In this case, the query can only return one row.) PL/pgSQL variables will be substituted into the query just as for commands that return no result, and the plan is cached in the same way. Also, the special variable `FOUND` is set to `true` if the query produced at least one row, or `false` if it produced no rows (see Section 43.5.5).

Note

One might expect that writing `SELECT` directly would accomplish this result, but at present the only accepted way to do it is `PERFORM`. A SQL command that can return rows, such as `SELECT`, will be rejected as an error unless it has an `INTO` clause as discussed in the next section.

An example:

```
PERFORM create_mv('cs_session_page_requests_mv', my_query);
```

### 43.5.3. Executing a Query with a Single-row Result

The result of a SQL command yielding a single row (possibly of multiple columns) can be assigned to a record variable, row-type variable, or list of scalar variables. This is done by writing the base SQL command and adding an `INTO` clause. For example,

```
SELECT select_expressions INTO [STRICT] target FROM ...;
INSERT ... RETURNING expressions INTO [STRICT] target;
UPDATE ... RETURNING expressions INTO [STRICT] target;
DELETE ... RETURNING expressions INTO [STRICT] target;
```

where `target` can be a record variable, a row variable, or a comma-separated list of simple variables and record/row fields. PL/pgSQL variables will be substituted into the rest of the query, and the plan is cached, just as described above for commands that do not return rows. This works for `SELECT`, `INSERT/UPDATE/DELETE` with `RETURNING`, and utility commands that return row-set results (such as `EXPLAIN`). Except for the `INTO` clause, the SQL command is the same as it would be written outside PL/pgSQL.

Tip

Note that this interpretation of `SELECT` with `INTO` is quite different from PostgreSQL's regular `SELECT INTO` command, wherein the `INTO` target is a newly created table. If you want to create a table from a `SELECT` result inside a PL/pgSQL function, use the syntax `CREATE TABLE ... AS SELECT`. 

1169
If a row or a variable list is used as target, the query's result columns must exactly match the structure of the target as to number and data types, or else a run-time error occurs. When a record variable is the target, it automatically configures itself to the row type of the query result columns.

The INTO clause can appear almost anywhere in the SQL command. Customarily it is written either just before or just after the list of select_expressions in a SELECT command, or at the end of the command for other command types. It is recommended that you follow this convention in case the PL/pgSQL parser becomes stricter in future versions.

If STRICT is not specified in the INTO clause, then target will be set to the first row returned by the query, or to nulls if the query returned no rows. (Note that “the first row” is not well-defined unless you’ve used ORDER BY.) Any result rows after the first row are discarded. You can check the special FOUND variable (see Section 43.5.5) to determine whether a row was returned:

```sql
SELECT * INTO myrec FROM emp WHERE empname = myname;
IF NOT FOUND THEN
    RAISE EXCEPTION 'employee % not found', myname;
END IF;
```

If the STRICT option is specified, the query must return exactly one row or a run-time error will be reported, either NO_DATA_FOUND (no rows) or TOO_MANY_ROWS (more than one row). You can use an exception block if you wish to catch the error, for example:

```sql
BEGIN
    SELECT * INTO STRICT myrec FROM emp WHERE empname = myname;
    EXCEPTION
        WHEN NO_DATA_FOUND THEN
            RAISE EXCEPTION 'employee % not found', myname;
        WHEN TOO_MANY_ROWS THEN
            RAISE EXCEPTION 'employee % not unique', myname;
END;
```

Successful execution of a command with STRICT always sets FOUND to true.

For INSERT/UPDATE/DELETE with RETURNING, PL/pgSQL reports an error for more than one returned row, even when STRICT is not specified. This is because there is no option such as ORDER BY with which to determine which affected row should be returned.

If print_strict_params is enabled for the function, then when an error is thrown because the requirements of STRICT are not met, the DETAIL part of the error message will include information about the parameters passed to the query. You can change the print_strict_params setting for all functions by setting plpgsql.print_strict_params, though only subsequent function compilations will be affected. You can also enable it on a per-function basis by using a compiler option, for example:

```sql
CREATE FUNCTION get_userid(username text) RETURNS int
AS $$
    #print_strict_params on
DECLARE
    userid int;
BEGIN
    SELECT users.userid INTO STRICT userid
    FROM users WHERE users.username = get_userid.username;
    RETURN userid;
END
$$ LANGUAGE plpgsql;
```
On failure, this function might produce an error message such as

ERROR:  query returned no rows
DETAIL:  parameters: $1 = 'nosuchuser'
CONTEXT:  PL/pgSQL function get_userid(text) line 6 at SQL statement

Note

The STRICT option matches the behavior of Oracle PL/SQL's SELECT INTO and related statements.

To handle cases where you need to process multiple result rows from a SQL query, see Section 43.6.6.

43.5.4. Executing Dynamic Commands

Oftentimes you will want to generate dynamic commands inside your PL/pgSQL functions, that is, commands that will involve different tables or different data types each time they are executed. PL/pgSQL's normal attempts to cache plans for commands (as discussed in Section 43.11.2) will not work in such scenarios. To handle this sort of problem, the EXECUTE statement is provided:

EXECUTE command-string [ INTO [STRICT] target ] [ USING expression [, ... ] ];

where command-string is an expression yielding a string (of type text) containing the command to be executed. The optional target is a record variable, a row variable, or a comma-separated list of simple variables and record/row fields, into which the results of the command will be stored. The optional USING expressions supply values to be inserted into the command.

No substitution of PL/pgSQL variables is done on the computed command string. Any required variable values must be inserted in the command string as it is constructed; or you can use parameters as described below.

Also, there is no plan caching for commands executed via EXECUTE. Instead, the command is always planned each time the statement is run. Thus the command string can be dynamically created within the function to perform actions on different tables and columns.

The INTO clause specifies where the results of a SQL command returning rows should be assigned. If a row or variable list is provided, it must exactly match the structure of the query's results (when a record variable is used, it will configure itself to match the result structure automatically). If multiple rows are returned, only the first will be assigned to the INTO variable. If no rows are returned, NULL is assigned to the INTO variable(s). If no INTO clause is specified, the query results are discarded.

If the STRICT option is given, an error is reported unless the query produces exactly one row.

The command string can use parameter values, which are referenced in the command as $1, $2, etc. These symbols refer to values supplied in the USING clause. This method is often preferable to inserting data values into the command string as text: it avoids run-time overhead of converting the values to text and back, and it is much less prone to SQL-injection attacks since there is no need for quoting or escaping. An example is:

EXECUTE 'SELECT count(*) FROM mytable WHERE inserted_by = $1 AND inserted <= $2'
           INTO c
USING checked_user, checked_date;

Note that parameter symbols can only be used for data values — if you want to use dynamically determined table or column names, you must insert them into the command string textually. For example, if the preceding query needed to be done against a dynamically selected table, you could do this:

```sql
EXECUTE 'SELECT count(*) FROM ' || quote_ident(tabname) || ' WHERE inserted_by = $1 AND inserted <= $2' INTO c
USING checked_user, checked_date;
```

A cleaner approach is to use `format()`’s `%I` specification for table or column names (strings separated by a newline are concatenated):

```sql
EXECUTE format('SELECT count(*) FROM %I ' || 'WHERE inserted_by = $1 AND inserted <= $2', tabname) INTO c
USING checked_user, checked_date;
```

Another restriction on parameter symbols is that they only work in `SELECT`, `INSERT`, `UPDATE`, and `DELETE` commands. In other statement types (generically called utility statements), you must insert values textually even if they are just data values.

An `EXECUTE` with a simple constant command string and some `USING` parameters, as in the first example above, is functionally equivalent to just writing the command directly in PL/pgSQL and allowing replacement of PL/pgSQL variables to happen automatically. The important difference is that `EXECUTE` will re-plan the command on each execution, generating a plan that is specific to the current parameter values; whereas PL/pgSQL may otherwise create a generic plan and cache it for reuse. In situations where the best plan depends strongly on the parameter values, it can be helpful to use `EXECUTE` to positively ensure that a generic plan is not selected.

```
SELECT INTO is not currently supported within EXECUTE; instead, execute a plain `SELECT` command and specify INTO as part of the EXECUTE itself.
```

---

### Note

The PL/pgSQL `EXECUTE` statement is not related to the `EXECUTE SQL` statement supported by the PostgreSQL server. The server’s `EXECUTE` statement cannot be used directly within PL/pgSQL functions (and is not needed).

---

### Example 43.1. Quoting Values In Dynamic Queries

When working with dynamic commands you will often have to handle escaping of single quotes. The recommended method for quoting fixed text in your function body is dollar quoting. (If you have legacy code that does not use dollar quoting, please refer to the overview in Section 43.12.1, which can save you some effort when translating said code to a more reasonable scheme.)

Dynamic values require careful handling since they might contain quote characters. An example using `format()` (this assumes that you are dollar quoting the function body so quote marks need not be doubled):

```
EXECUTE format('UPDATE tbl SET %I = $1 ' || 'WHERE key = $2', colname) USING newvalue, keyvalue;
```
It is also possible to call the quoting functions directly:

```
EXECUTE 'UPDATE tbl SET 
    | quote_ident(colname) 
    | ' = ' 
    | quote_literal(newvalue) 
    | ' WHERE key = ' 
    | quote_literal(keyvalue);
```

This example demonstrates the use of the `quote_ident` and `quote_literal` functions (see Section 9.4). For safety, expressions containing column or table identifiers should be passed through `quote_ident` before insertion in a dynamic query. Expressions containing values that should be literal strings in the constructed command should be passed through `quote_literal`. These functions take the appropriate steps to return the input text enclosed in double or single quotes respectively, with any embedded special characters properly escaped.

Because `quote_literal` is labeled `STRICT`, it will always return null when called with a null argument. In the above example, if `newvalue` or `keyvalue` were null, the entire dynamic query string would become null, leading to an error from `EXECUTE`. You can avoid this problem by using the `quote_nullable` function, which works the same as `quote_literal` except that when called with a null argument it returns the string `NULL`. For example,

```
EXECUTE 'UPDATE tbl SET 
    | quote_ident(colname) 
    | ' = ' 
    | quote_nullable(newvalue) 
    | ' WHERE key = ' 
    | quote_nullable(keyvalue);
```

If you are dealing with values that might be null, you should usually use `quote_nullable` in place of `quote_literal`.

As always, care must be taken to ensure that null values in a query do not deliver unintended results. For example the `WHERE` clause

```
'WHERE key = ' | quote_nullable(keyvalue)
```

will never succeed if `keyvalue` is null, because the result of using the equality operator `=` with a null operand is always null. If you wish null to work like an ordinary key value, you would need to rewrite the above as

```
'WHERE key IS NOT DISTINCT FROM ' | quote_nullable(keyvalue)
```

(At present, `IS NOT DISTINCT FROM` is handled much less efficiently than `=`, so don’t do this unless you must. See Section 9.2 for more information on nulls and `IS DISTINCT`.)

Note that dollar quoting is only useful for quoting fixed text. It would be a very bad idea to try to write this example as:

```
EXECUTE 'UPDATE tbl SET 
    | quote_ident(colname) 
    | ' = $$' 
    | newvalue 
    | $$ WHERE key = ' 
    | quote_literal(keyvalue);
```
because it would break if the contents of \texttt{newvalue} happened to contain $\$$. The same objection would apply to any other dollar-quoting delimiter you might pick. So, to safely quote text that is not known in advance, you \textit{must} use \texttt{quote\_literal}, \texttt{quote\_nullable}, or \texttt{quote\_ident}, as appropriate.

Dynamic SQL statements can also be safely constructed using the \texttt{format} function (see Section 9.4). For example:

\begin{verbatim}
EXECUTE format('UPDATE tbl SET %I = %L '
  'WHERE key = %L', colname, newvalue, keyvalue);
\end{verbatim}

\%I is equivalent to \texttt{quote\_ident}, and \%L is equivalent to \texttt{quote\_nullable}. The \texttt{format} function can be used in conjunction with the \texttt{USING} clause:

\begin{verbatim}
EXECUTE format('UPDATE tbl SET %I = $1 WHERE key = $2', colname)
  USING newvalue, keyvalue;
\end{verbatim}

This form is better because the variables are handled in their native data type format, rather than unconditionally converting them to text and quoting them via \%L. It is also more efficient.

A much larger example of a dynamic command and \texttt{EXECUTE} can be seen in Example 43.10, which builds and executes a \texttt{CREATE\_FUNCTION} command to define a new function.

\section*{43.5.5. Obtaining the Result Status}

There are several ways to determine the effect of a command. The first method is to use the \texttt{GET DIAGNOSTICS} command, which has the form:

\begin{verbatim}
GET [ CURRENT ] DIAGNOSTICS variable { = | := } item [...];
\end{verbatim}

This command allows retrieval of system status indicators. \texttt{CURRENT} is a noise word (but see also \texttt{GET STACKED DIAGNOSTICS} in Section 43.6.8.1). Each \texttt{item} is a key word identifying a status value to be assigned to the specified \texttt{variable} (which should be of the right data type to receive it). The currently available status items are shown in Table 43.1. Colon-equal (\:=) can be used instead of the SQL-standard = token. An example:

\begin{verbatim}
GET DIAGNOSTICS integer_var = ROW_COUNT;
\end{verbatim}

\begin{table}[h]
\centering
\caption{Available Diagnostics Items}
\begin{tabular}{|l|l|l|}
\hline
Name & Type & Description \\
\hline
ROW\_COUNT & bigint & the number of rows processed by the most recent SQL command \\
\hline
RESULT\_OID & oid & the OID of the last row inserted by the most recent SQL command (only useful after an \texttt{INSERT} command into a table having OIDs) \\
\hline
PG\_CONTEXT & text & line(s) of text describing the current call stack (see Section 43.6.9) \\
\hline
\end{tabular}
\end{table}

The second method to determine the effects of a command is to check the special variable named \texttt{FOUND}, which is of type \texttt{boolean}. \texttt{FOUND} starts out false within each PL/pgSQL function call. It is set by each of the following types of statements:
• A `SELECT INTO` statement sets `FOUND` true if a row is assigned, false if no row is returned.

• A `PERFORM` statement sets `FOUND` true if it produces (and discards) one or more rows, false if no row is produced.

• `UPDATE`, `INSERT`, and `DELETE` statements set `FOUND` true if at least one row is affected, false if no row is affected.

• A `FETCH` statement sets `FOUND` true if it returns a row, false if no row is returned.

• A `MOVE` statement sets `FOUND` true if it successfully repositions the cursor, false otherwise.

• A `FOR` or `FOREACH` statement sets `FOUND` true if it iterates one or more times, else false. `FOUND` is set this way when the loop exits; inside the execution of the loop, `FOUND` is not modified by the loop statement, although it might be changed by the execution of other statements within the loop body.

• `RETURN QUERY` and `RETURN QUERY EXECUTE` statements set `FOUND` true if the query returns at least one row, false if no row is returned.

Other PL/pgSQL statements do not change the state of `FOUND`. Note in particular that `EXECUTE` changes the output of `GET DIAGNOSTICS`, but does not change `FOUND`.

`FOUND` is a local variable within each PL/pgSQL function; any changes to it affect only the current function.

### 43.5.6. Doing Nothing At All

Sometimes a placeholder statement that does nothing is useful. For example, it can indicate that one arm of an `if/then/else` chain is deliberately empty. For this purpose, use the `NULL` statement:

```
NULL;
```

For example, the following two fragments of code are equivalent:

```
BEGIN
  y := x / 0;
EXCEPTION
  WHEN division_by_zero THEN
    NULL; -- ignore the error
END;

BEGIN
  y := x / 0;
EXCEPTION
  WHEN division_by_zero THEN  -- ignore the error
END;
```

Which is preferable is a matter of taste.

---

**Note**

In Oracle's PL/SQL, empty statement lists are not allowed, and so `NULL` statements are *required* for situations such as this. PL/pgSQL allows you to just write nothing, instead.
43.6. Control Structures

Control structures are probably the most useful (and important) part of PL/pgSQL. With PL/pgSQL’s control structures, you can manipulate PostgreSQL data in a very flexible and powerful way.

43.6.1. Returning From a Function

There are two commands available that allow you to return data from a function: RETURN and RETURN NEXT.

43.6.1.1. RETURN

```
RETURN expression;
```

RETURN with an expression terminates the function and returns the value of expression to the caller. This form is used for PL/pgSQL functions that do not return a set.

In a function that returns a scalar type, the expression’s result will automatically be cast into the function’s return type as described for assignments. But to return a composite (row) value, you must write an expression delivering exactly the requested column set. This may require use of explicit casting.

If you declared the function with output parameters, write just RETURN with no expression. The current values of the output parameter variables will be returned.

If you declared the function to return void, a RETURN statement can be used to exit the function early; but do not write an expression following RETURN.

The return value of a function cannot be left undefined. If control reaches the end of the top-level block of the function without hitting a RETURN statement, a run-time error will occur. This restriction does not apply to functions with output parameters and functions returning void, however. In those cases a RETURN statement is automatically executed if the top-level block finishes.

Some examples:

```sql
-- functions returning a scalar type
RETURN 1 + 2;
RETURN scalar_var;

-- functions returning a composite type
RETURN composite_type_var;
RETURN (1, 2, 'three':text);  -- must cast columns to correct types
```

43.6.1.2. RETURN NEXT and RETURN QUERY

```
RETURN NEXT expression;
RETURN QUERY query;
RETURN QUERY EXECUTE command-string [ USING expression [, ... ] ];
```

When a PL/pgSQL function is declared to return SETOF sometype, the procedure to follow is slightly different. In that case, the individual items to return are specified by a sequence of RETURN NEXT or RETURN QUERY commands, and then a final RETURN command with no argument is used to indicate that the function has finished executing. RETURN NEXT can be used with both scalar and composite data types; with a composite result type, an entire “table” of results will be returned.
RETURN QUERY appends the results of executing a query to the function’s result set. RETURN NEXT and RETURN QUERY can be freely intermixed in a single set-returning function, in which case their results will be concatenated.

RETURN NEXT and RETURN QUERY do not actually return from the function — they simply append zero or more rows to the function’s result set. Execution then continues with the next statement in the PL/pgSQL function. As successive RETURN NEXT or RETURN QUERY commands are executed, the result set is built up. A final RETURN, which should have no argument, causes control to exit the function (or you can just let control reach the end of the function).

RETURN QUERY has a variant RETURN QUERY EXECUTE, which specifies the query to be executed dynamically. Parameter expressions can be inserted into the computed query string via USING, in just the same way as in the EXECUTE command.

If you declared the function with output parameters, write just RETURN NEXT with no expression. On each execution, the current values of the output parameter variable(s) will be saved for eventual return as a row of the result. Note that you must declare the function as returning SETOF record when there are multiple output parameters, or SETOF sometype when there is just one output parameter of type sometype, in order to create a set-returning function with output parameters.

Here is an example of a function using RETURN NEXT:

```
CREATE TABLE foo (fooid INT, foosubid INT, fooname TEXT);
INSERT INTO foo VALUES (1, 2, 'three');
INSERT INTO foo VALUES (4, 5, 'six');

CREATE OR REPLACE FUNCTION get_all_foo() RETURNS SETOF foo AS
$BODY$
DECLARE
  r foo%rowtype;
BEGIN
  FOR r IN
    SELECT * FROM foo WHERE fooid > 0
  LOOP
    -- can do some processing here
    RETURN NEXT r; -- return current row of SELECT
  END LOOP;
RETURN;
$BODY$
LANGUAGE plpgsql;
SELECT * FROM get_all_foo();
```

Here is an example of a function using RETURN QUERY:

```
CREATE FUNCTION get_available_flightid(date) RETURNS SETOF integer AS
$BODY$
BEGIN
  RETURN QUERY SELECT flightid 
  FROM flight 
  WHERE flightdate >= $1 
  AND flightdate < ($1 + 1); 

  -- Since execution is not finished, we can check whether rows were returned
$BODY$
LANGUAGE plpgsql;
SELECT * FROM get_available_flightid('2023-01-01');
```
-- and raise exception if not.
IF NOT FOUND THEN
  RAISE EXCEPTION 'No flight at %.', $1;
END IF;

RETURN;
END
$BODY$
LANGUAGE plpgsql;

-- Returns available flights or raises exception if there are no
-- available flights.
SELECT * FROM get_available_flightid(CURRENT_DATE);

Note
The current implementation of RETURN NEXT and RETURN QUERY stores the entire result set before returning from the function, as discussed above. That means that if a PL/pgSQL function produces a very large result set, performance might be poor: data will be written to disk to avoid memory exhaustion, but the function itself will not return until the entire result set has been generated. A future version of PL/pgSQL might allow users to define set-returning functions that do not have this limitation. Currently, the point at which data begins being written to disk is controlled by the work_mem configuration variable. Administrators who have sufficient memory to store larger result sets in memory should consider increasing this parameter.

43.6.2. Returning From a Procedure

A procedure does not have a return value. A procedure can therefore end without a RETURN statement. If you wish to use a RETURN statement to exit the code early, write just RETURN with no expression.

If the procedure has output parameters, the final values of the output parameter variables will be returned to the caller.

43.6.3. Calling a Procedure

A PL/pgSQL function, procedure, or DO block can call a procedure using CALL. Output parameters are handled differently from the way that CALL works in plain SQL. Each INOUT parameter of the procedure must correspond to a variable in the CALL statement, and whatever the procedure returns is assigned back to that variable after it returns. For example:

CREATE PROCEDURE triple(INOUT x int)
LANGUAGE plpgsql
AS $$
BEGIN
  x := x * 3;
END;
$$;

DO $$
DECLARE myvar int := 5;
BEGIN
  CALL triple(myvar);
  RAISE NOTICE 'myvar = %', myvar; -- prints 15
END
$$;
$\$;

## 43.6.4. Conditionals

IF and CASE statements let you execute alternative commands based on certain conditions. PL/pgSQL has three forms of IF:

- IF ... THEN ... END IF
- IF ... THEN ... ELSE ... END IF
- IF ... THEN ... ELSIF ... THEN ... ELSE ... END IF

and two forms of CASE:

- CASE ... WHEN ... THEN ... ELSE ... END CASE
- CASE WHEN ... THEN ... ELSE ... END CASE

### 43.6.4.1. IF–THEN

\[
\text{IF boolean-expression THEN} \\
\hspace{1cm} \text{statements} \\
\hspace{1cm} \text{END IF;}
\]

IF–THEN statements are the simplest form of IF. The statements between THEN and END IF will be executed if the condition is true. Otherwise, they are skipped.

Example:

```sql
IF v_user_id <> 0 THEN 
    UPDATE users SET email = v_email WHERE user_id = v_user_id;
END IF;
```

### 43.6.4.2. IF–THEN–ELSE

\[
\text{IF boolean-expression THEN} \\
\hspace{1cm} \text{statements} \\
\hspace{1cm} \text{ELSE} \\
\hspace{1cm} \hspace{1cm} \text{statements} \\
\hspace{1cm} \text{END IF;}
\]

IF–THEN–ELSE statements add to IF–THEN by letting you specify an alternative set of statements that should be executed if the condition is not true. (Note this includes the case where the condition evaluates to NULL.)

Examples:

```sql
IF parentid IS NULL OR parentid = '' 
THEN  
    RETURN fullname;
ELSE 
    RETURN hp_true_filename(parentid) || '/' || fullname;
END IF;
```
IF v_count > 0 THEN 
    INSERT INTO users_count (count) VALUES (v_count);
    RETURN 't';
ELSE 
    RETURN 'f';
END IF;

43.6.4.3. **IF–THEN–ELSIF**

IF boolean-expression THEN  
    statements  
[ ELSEIF boolean-expression THEN  
    statements  
[ ELSEIF boolean-expression THEN  
    statements  
    ...  
] ] ] [ ELSE  
    statements ]
END IF;

Sometimes there are more than just two alternatives. **IF–THEN–ELSIF** provides a convenient method of checking several alternatives in turn. The IF conditions are tested successively until the first one that is true is found. Then the associated statement(s) are executed, after which control passes to the next statement after END IF. (Any subsequent IF conditions are not tested.) If none of the IF conditions is true, then the ELSE block (if any) is executed.

Here is an example:

IF number = 0 THEN 
    result := 'zero';
ELSIF number > 0 THEN 
    result := 'positive';
ELSIF number < 0 THEN 
    result := 'negative';
ELSE 
    -- hmm, the only other possibility is that number is null
    result := 'NULL';
END IF;

The key word ELSIF can also be spelled ELSEIF.

An alternative way of accomplishing the same task is to nest IF–THEN–ELSE statements, as in the following example:

IF demo_row.sex = 'm' THEN 
    pretty_sex := 'man';
ELSE 
    IF demo_row.sex = 'f' THEN 
        pretty Sex := 'woman';
    END IF;
END IF;

However, this method requires writing a matching END IF for each IF, so it is much more cumbersome than using ELSIF when there are many alternatives.
43.6.4.4. Simple CASE

```plpgsql
case search-expression
    when expression [ , expression [ ... ] ] then
      statements
      [ when expression [ , expression [ ... ] ] then
      statements
      ... ]
      [ else
      statements ]
end case;
```

The simple form of CASE provides conditional execution based on equality of operands. The `search-expression` is evaluated (once) and successively compared to each `expression` in the WHEN clauses. If a match is found, then the corresponding `statements` are executed, and then control passes to the next statement after END CASE. (Subsequent WHEN expressions are not evaluated.) If no match is found, the ELSE statements are executed; but if ELSE is not present, then a CASE NOT FOUND exception is raised.

Here is a simple example:

```plpgsql
case x
  when 1, 2 then
    msg := 'one or two';
  else
    msg := 'other value than one or two';
end case;
```

43.6.4.5. Searched CASE

```plpgsql
case
  when boolean-expression then
    statements
    [ when boolean-expression then
    statements
    ... ]
    [ else
    statements ]
end case;
```

The searched form of CASE provides conditional execution based on truth of Boolean expressions. Each WHEN clause’s `boolean-expression` is evaluated in turn, until one is found that yields true. Then the corresponding `statements` are executed, and then control passes to the next statement after END CASE. (Subsequent WHEN expressions are not evaluated.) If no true result is found, the ELSE statements are executed; but if ELSE is not present, then a CASE NOT FOUND exception is raised.

Here is an example:

```plpgsql
case
  when x between 0 and 10 then
    msg := 'value is between zero and ten';
  when x between 11 and 20 then
    msg := 'value is between eleven and twenty';
end case;
```
This form of CASE is entirely equivalent to IF-THEN-ELSIF, except for the rule that reaching an omitted ELSE clause results in an error rather than doing nothing.

### 43.6.5. Simple Loops

With the LOOP, EXIT, CONTINUE, WHILE, FOR, and FOREACH statements, you can arrange for your PL/pgSQL function to repeat a series of commands.

#### 43.6.5.1. LOOP

```plaintext
[ <<label>> ]
LOOP
  statements
END LOOP [ label ];
```

LOOP defines an unconditional loop that is repeated indefinitely until terminated by an EXIT or RETURN statement. The optional label can be used by EXIT and CONTINUE statements within nested loops to specify which loop those statements refer to.

#### 43.6.5.2. EXIT

```plaintext
EXIT [ label ] [ WHEN boolean-expression ];
```

If no label is given, the innermost loop is terminated and the statement following END LOOP is executed next. If label is given, it must be the label of the current or some outer level of nested loop or block. Then the named loop or block is terminated and control continues with the statement after the loop's/block's corresponding END.

If WHEN is specified, the loop exit occurs only if boolean-expression is true. Otherwise, control passes to the statement after EXIT.

EXIT can be used with all types of loops; it is not limited to use with unconditional loops.

When used with a BEGIN block, EXIT passes control to the next statement after the end of the block. Note that a label must be used for this purpose; an unlabeled EXIT is never considered to match a BEGIN block. (This is a change from pre-8.4 releases of PostgreSQL, which would allow an unlabeled EXIT to match a BEGIN block.)

Examples:

```plaintext
LOOP
  -- some computations
  IF count > 0 THEN
    EXIT;  -- exit loop
  END IF;
END LOOP;

LOOP
  -- some computations
  EXIT WHEN count > 0;  -- same result as previous example
END LOOP;

<<ablock>>
BEGIN
  -- some computations
  IF stocks > 100000 THEN
```
EXIT ablock;  -- causes exit from the BEGIN block
END IF;
    -- computations here will be skipped when stocks > 100000
END;

43.6.5.3. CONTINUE

CONTINUE [ label ] [ WHEN boolean-expression ];

If no label is given, the next iteration of the innermost loop is begun. That is, all statements remaining
in the loop body are skipped, and control returns to the loop control expression (if any) to determine
whether another loop iteration is needed. If label is present, it specifies the label of the loop whose
execution will be continued.

If WHEN is specified, the next iteration of the loop is begun only if boolean-expression is true.
Otherwise, control passes to the statement after CONTINUE.

CONTINUE can be used with all types of loops; it is not limited to use with unconditional loops.

Examples:

LOOP
    -- some computations
    EXIT WHEN count > 100;
    CONTINUE WHEN count < 50;
    -- some computations for count IN [50 .. 100]
END LOOP;

43.6.5.4. WHILE

[ <<label>> ]
WHILE boolean-expression LOOP
    statements
END LOOP [ label ];

The WHILE statement repeats a sequence of statements so long as the boolean-expression
evaluates to true. The expression is checked just before each entry to the loop body.

For example:

WHILE amount_owed > 0 AND gift_certificate_balance > 0 LOOP
    -- some computations here
END LOOP;

WHILE NOT done LOOP
    -- some computations here
END LOOP;

43.6.5.5. FOR (Integer Variant)

[ <<label>> ]
FOR name IN [ REVERSE ] expression .. expression [ BY expression ]
    LOOP
        statements
    END LOOP [ label ];
This form of FOR creates a loop that iterates over a range of integer values. The variable name is automatically defined as type integer and exists only inside the loop (any existing definition of the variable name is ignored within the loop). The two expressions giving the lower and upper bound of the range are evaluated once when entering the loop. If the BY clause isn't specified the iteration step is 1, otherwise it's the value specified in the BY clause, which again is evaluated once on loop entry. If REVERSE is specified then the step value is subtracted, rather than added, after each iteration.

Some examples of integer FOR loops:

```
FOR i IN 1..10 LOOP
    -- i will take on the values 1,2,3,4,5,6,7,8,9,10 within the loop
END LOOP;
```

```
FOR i IN REVERSE 10..1 LOOP
    -- i will take on the values 10,9,8,7,6,5,4,3,2,1 within the loop
END LOOP;
```

```
FOR i IN REVERSE 10..1 BY 2 LOOP
    -- i will take on the values 10,8,6,4,2 within the loop
END LOOP;
```

If the lower bound is greater than the upper bound (or less than, in the REVERSE case), the loop body is not executed at all. No error is raised.

If a label is attached to the FOR loop then the integer loop variable can be referenced with a qualified name, using that label.

### 43.6.6. Looping Through Query Results

Using a different type of FOR loop, you can iterate through the results of a query and manipulate that data accordingly. The syntax is:

```
[ <<label>> ]
FOR target IN query LOOP
    statements
END LOOP [ label ];
```

The target is a record variable, row variable, or comma-separated list of scalar variables. The target is successively assigned each row resulting from the query and the loop body is executed for each row. Here is an example:

```
CREATE FUNCTION cs_refresh_mviews() RETURNS integer AS $$
DECLARE
    mviews RECORD;
BEGIN
    RAISE NOTICE 'Refreshing materialized views...';

    FOR mviews IN SELECT * FROM cs_materialized_views ORDER BY sort_key LOOP
        -- Now "mviews" has one record from cs_materialized_views
        RAISE NOTICE 'Refreshing materialized view %s ...', quote_ident(mviews.mv_name);
    END LOOP;

    RAISE NOTICE 'Finished refreshing materialized views...';
END $$ LANGUAGE plpgsql;
```

```
EXECUTE format('TRUNCATE TABLE %I', mviews.mv_name);
EXECUTE format('INSERT INTO %I %s', mviews.mv_name, mviews.mv_query);
END LOOP;

RAISE NOTICE 'Done refreshing materialized views.';
RETURN 1;
END;
$$ LANGUAGE plpgsql;

If the loop is terminated by an EXIT statement, the last assigned row value is still accessible after the loop.

The query used in this type of FOR statement can be any SQL command that returns rows to the caller: SELECT is the most common case, but you can also use INSERT, UPDATE, or DELETE with a RETURNING clause. Some utility commands such as EXPLAIN will work too.

PL/pgSQL variables are substituted into the query text, and the query plan is cached for possible re-use, as discussed in detail in Section 43.11.1 and Section 43.11.2.

The FOR-IN-EXECUTE statement is another way to iterate over rows:

```
[[label]]
FOR target IN EXECUTE text_expression [ USING expression [, ... ] ]
LOOP
  statements
END LOOP [ label ];
```

This is like the previous form, except that the source query is specified as a string expression, which is evaluated and replanned on each entry to the FOR loop. This allows the programmer to choose the speed of a preplanned query or the flexibility of a dynamic query, just as with a plain EXECUTE statement. As with EXECUTE, parameter values can be inserted into the dynamic command via USING.

Another way to specify the query whose results should be iterated through is to declare it as a cursor. This is described in Section 43.7.4.

### 43.6.7. Looping Through Arrays

The FOREACH loop is much like a FOR loop, but instead of iterating through the rows returned by a SQL query, it iterates through the elements of an array value. (In general, FOREACH is meant for looping through components of a composite-valued expression; variants for looping through composites besides arrays may be added in future.) The FOREACH statement to loop over an array is:

```
[[label]]
FOREACH target [ SLICE number ] IN ARRAY expression LOOP
  statements
END LOOP [ label ];
```

Without SLICE, or if SLICE 0 is specified, the loop iterates through individual elements of the array produced by evaluating the expression. The target variable is assigned each element value in sequence, and the loop body is executed for each element. Here is an example of looping through the elements of an integer array:

```
CREATE FUNCTION sum(int[]) RETURNS int8 AS $$
DECLARE
  s int8 := 0;
  x int;
```

BEGIN
    FOREACH x IN ARRAY $1
    LOOP
        s := s + x;
    END LOOP;
    RETURN s;
END;
$$ LANGUAGE plpgsql;

The elements are visited in storage order, regardless of the number of array dimensions. Although the target is usually just a single variable, it can be a list of variables when looping through an array of composite values (records). In that case, for each array element, the variables are assigned from successive columns of the composite value.

With a positive SLICE value, FOREACH iterates through slices of the array rather than single elements. The SLICE value must be an integer constant not larger than the number of dimensions of the array. The target variable must be an array, and it receives successive slices of the array value, where each slice is of the number of dimensions specified by SLICE. Here is an example of iterating through one-dimensional slices:

CREATE FUNCTION scan_rows(int[]) RETURNS void AS $$
DECLARE
    x int[];
BEGIN
    FOREACH x SLICE 1 IN ARRAY $1
    LOOP
        RAISE NOTICE 'row = %', x;
    END LOOP;
END;
$$ LANGUAGE plpgsql;

SELECT scan_rows(ARRAY[[1,2,3],[4,5,6],[7,8,9],[10,11,12]]);

NOTICE:  row = {1,2,3}
NOTICE:  row = {4,5,6}
NOTICE:  row = {7,8,9}
NOTICE:  row = {10,11,12}

43.6.8. Trapping Errors

By default, any error occurring in a PL/pgSQL function aborts execution of the function, and indeed of the surrounding transaction as well. You can trap errors and recover from them by using a BEGIN block with an EXCEPTION clause. The syntax is an extension of the normal syntax for a BEGIN block:

[ <<label>> ]
[ DECLARE
    declarations ]
BEGIN
    statements
EXCEPTION
    WHEN condition [ OR condition ... ] THEN
        handler_statements
    [ WHEN condition [ OR condition ... ] THEN
        handler_statements
    ... ]
END;
If no error occurs, this form of block simply executes all the statements, and then control passes to the next statement after END. But if an error occurs within the statements, further processing of the statements is abandoned, and control passes to the EXCEPTION list. The list is searched for the first condition matching the error that occurred. If a match is found, the corresponding handler_statements are executed, and then control passes to the next statement after END. If no match is found, the error propagates out as though the EXCEPTION clause were not there at all: the error can be caught by an enclosing block with EXCEPTION, or if there is none it aborts processing of the function.

The condition names can be any of those shown in Appendix A. A category name matches any error within its category. The special condition name OTHERS matches every error type except QUERY_CANCELED and ASSERT_FAILURE. (It is possible, but often unwise, to trap those two error types by name.) Condition names are not case-sensitive. Also, an error condition can be specified by SQLSTATE code; for example these are equivalent:

WHEN division_by_zero THEN ...  
WHEN SQLSTATE '22012' THEN ...

If a new error occurs within the selected handler_statements, it cannot be caught by this EXCEPTION clause, but is propagated out. A surrounding EXCEPTION clause could catch it.

When an error is caught by an EXCEPTION clause, the local variables of the PL/pgSQL function remain as they were when the error occurred, but all changes to persistent database state within the block are rolled back. As an example, consider this fragment:

```
INSERT INTO mytab(firstname, lastname) VALUES('Tom', 'Jones');  
BEGIN  
  UPDATE mytab SET firstname = 'Joe' WHERE lastname = 'Jones';  
x := x + 1;  
y := x / 0;  
EXCEPTION  
  WHEN division_by_zero THEN  
    RAISE NOTICE 'caught division_by_zero';  
    RETURN x;  
END;
```

When control reaches the assignment to y, it will fail with a division_by_zero error. This will be caught by the EXCEPTION clause. The value returned in the RETURN statement will be the incremented value of x, but the effects of the UPDATE command will have been rolled back. The INSERT command preceding the block is not rolled back, however, so the end result is that the database contains Tom Jones not Joe Jones.

Tip

A block containing an EXCEPTION clause is significantly more expensive to enter and exit than a block without one. Therefore, don’t use EXCEPTION without need.

Example 43.2. Exceptions with UPDATE/INSERT

This example uses exception handling to perform either UPDATE or INSERT, as appropriate. It is recommended that applications use INSERT with ON CONFLICT DO UPDATE rather than actually using this pattern. This example serves primarily to illustrate use of PL/pgSQL control flow structures:

```
CREATE TABLE db (a INT PRIMARY KEY, b TEXT);
```
CREATE FUNCTION merge_db(key INT, data TEXT) RETURNS VOID AS $$
BEGIN
  LOOP
    -- first try to update the key
    UPDATE db SET b = data WHERE a = key;
    IF found THEN
      RETURN;
    END IF;
    -- not there, so try to insert the key
    -- if someone else inserts the same key concurrently,
    -- we could get a unique-key failure
    BEGIN
      INSERT INTO db(a,b) VALUES (key, data);
      RETURN;
    EXCEPTION WHEN unique_violation THEN
      -- Do nothing, and loop to try the UPDATE again.
    END;
  END LOOP;
END;
$$
LANGUAGE plpgsql;

SELECT merge_db(1, 'david');
SELECT merge_db(1, 'dennis');

This coding assumes the unique_violation error is caused by the INSERT, and not by, say, an INSERT in a trigger function on the table. It might also misbehave if there is more than one unique index on the table, since it will retry the operation regardless of which index caused the error. More safety could be had by using the features discussed next to check that the trapped error was the one expected.

43.6.8.1. Obtaining Information About an Error

Exception handlers frequently need to identify the specific error that occurred. There are two ways to get information about the current exception in PL/pgSQL: special variables and the GET STACKED DIAGNOSTICS command.

Within an exception handler, the special variable SQLSTATE contains the error code that corresponds to the exception that was raised (refer to Table A.1 for a list of possible error codes). The special variable SQLERRM contains the error message associated with the exception. These variables are undefined outside exception handlers.

Within an exception handler, one may also retrieve information about the current exception by using the GET STACKED DIAGNOSTICS command, which has the form:

GET STACKED DIAGNOSTICS variable { = | := } item [ , ... ];

Each item is a key word identifying a status value to be assigned to the specified variable (which should be of the right data type to receive it). The currently available status items are shown in Table 43.2.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>RETURNED_SQLSTATE</td>
<td>text</td>
<td>the SQLSTATE error code of the exception</td>
</tr>
</tbody>
</table>

1188
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>COLUMN_NAME</td>
<td>text</td>
<td>the name of the column related to exception</td>
</tr>
<tr>
<td>CONSTRAINT_NAME</td>
<td>text</td>
<td>the name of the constraint related to exception</td>
</tr>
<tr>
<td>PG_DATATYPE_NAME</td>
<td>text</td>
<td>the name of the data type related to exception</td>
</tr>
<tr>
<td>MESSAGE_TEXT</td>
<td>text</td>
<td>the text of the exception's primary message</td>
</tr>
<tr>
<td>TABLE_NAME</td>
<td>text</td>
<td>the name of the table related to exception</td>
</tr>
<tr>
<td>SCHEMA_NAME</td>
<td>text</td>
<td>the name of the schema related to exception</td>
</tr>
<tr>
<td>PG_EXCEPTION_DETAIL</td>
<td>text</td>
<td>the text of the exception's detail message, if any</td>
</tr>
<tr>
<td>PG_EXCEPTION_HINT</td>
<td>text</td>
<td>the text of the exception's hint message, if any</td>
</tr>
<tr>
<td>PG_EXCEPTION_CONTEXT</td>
<td>text</td>
<td>line(s) of text describing the call stack at the time of the exception</td>
</tr>
</tbody>
</table>

If the exception did not set a value for an item, an empty string will be returned.

Here is an example:

```sql
DECLARE
  text_var1 text;
  text_var2 text;
  text_var3 text;
BEGIN
  -- some processing which might cause an exception
  ...
EXCEPTION WHEN OTHERS THEN
  GET STACKED DIAGNOSTICS text_var1 = MESSAGE_TEXT,
                          text_var2 = PG_EXCEPTION_DETAIL,
                          text_var3 = PG_EXCEPTION_HINT;
END;
```

### 43.6.9. Obtaining Execution Location Information

The **GET DIAGNOSTICS** command, previously described in Section 43.5.5, retrieves information about current execution state (whereas the **GET STACKED DIAGNOSTICS** command discussed above reports information about the execution state as of a previous error). Its **PG_CONTEXT** status item is useful for identifying the current execution location. **PG_CONTEXT** returns a text string with line(s) of text describing the call stack. The first line refers to the current function and currently executing **GET DIAGNOSTICS** command. The second and any subsequent lines refer to calling functions further up the call stack. For example:

```sql
CREATE OR REPLACE FUNCTION outer_func() RETURNS integer AS $$
BEGIN
  RETURN inner_func();
END;
$$ LANGUAGE plpgsql;
```
CREATE OR REPLACE FUNCTION inner_func() RETURNS integer AS $$
DECLARE
  stack text;
BEGIN
  GET DIAGNOSTICS stack = PG_CONTEXT;
  RAISE NOTICE E'--- Call Stack ---
%', stack;
  RETURN 1;
END;
$$ LANGUAGE plpgsql;

SELECT outer_func();

NOTICE: --- Call Stack ---
PL/pgSQL function inner_func() line 5 at GET DIAGNOSTICS
PL/pgSQL function outer_func() line 3 at RETURN
CONTEXT:  PL/pgSQL function outer_func() line 3 at RETURN
outer_func
------------
1
(1 row)

GET STACKED DIAGNOSTICS ... PG_EXCEPTION_CONTEXT returns the same sort of stack trace, but describing the location at which an error was detected, rather than the current location.

43.7. Cursors

Rather than executing a whole query at once, it is possible to set up a cursor that encapsulates the query, and then read the query result a few rows at a time. One reason for doing this is to avoid memory overrun when the result contains a large number of rows. (However, PL/pgSQL users do not normally need to worry about that, since FOR loops automatically use a cursor internally to avoid memory problems.) A more interesting usage is to return a reference to a cursor that a function has created, allowing the caller to read the rows. This provides an efficient way to return large row sets from functions.

43.7.1. Declaring Cursor Variables

All access to cursors in PL/pgSQL goes through cursor variables, which are always of the special data type refcursor. One way to create a cursor variable is just to declare it as a variable of type refcursor. Another way is to use the cursor declaration syntax, which in general is:

```
name [ [ NO ] SCROLL ] CURSOR [ ( arguments ) ] FOR query;
```

(FOR can be replaced by IS for Oracle compatibility.) If SCROLL is specified, the cursor will be capable of scrolling backward; if NO SCROLL is specified, backward fetches will be rejected; if neither specification appears, it is query-dependent whether backward fetches will be allowed. arguments, if specified, is a comma-separated list of pairs name datatype that define names to be replaced by parameter values in the given query. The actual values to substitute for these names will be specified later, when the cursor is opened.

Some examples:

```
DECLARE
  curs1 refcursor;
  curs2 CURSOR FOR SELECT * FROM tenk1;
```
curs3 CURSOR (key integer) FOR SELECT * FROM tenk1 WHERE unique1 = key;

All three of these variables have the data type refcursor, but the first can be used with any query, while the second has a fully specified query already bound to it, and the last has a parameterized query bound to it. \(key\) will be replaced by an integer parameter value when the cursor is opened.) The variable curs1 is said to be unbound since it is not bound to any particular query.

### 43.7.2. Opening Cursors

Before a cursor can be used to retrieve rows, it must be opened. (This is the equivalent action to the SQL command `DECLARE CURSOR`.) PL/pgSQL has three forms of the `OPEN` statement, two of which use unbound cursor variables while the third uses a bound cursor variable.

#### Note

Bound cursor variables can also be used without explicitly opening the cursor, via the `FOR` statement described in Section 43.7.4.

### 43.7.2.1. `OPEN` FOR `query`

```sql
OPEN unbound_cursorvar [ [ NO ] SCROLL ] FOR query;
```

The cursor variable is opened and given the specified query to execute. The cursor cannot be open already, and it must have been declared as an unbound cursor variable (that is, as a simple refcursor variable). The query must be a SELECT, or something else that returns rows (such as `EXPLAIN`). The query is treated in the same way as other SQL commands in PL/pgSQL: PL/pgSQL variable names are substituted, and the query plan is cached for possible reuse. When a PL/pgSQL variable is substituted into the cursor query, the value that is substituted is the one it has at the time of the `OPEN`; subsequent changes to the variable will not affect the cursor’s behavior. The SCROLL and NO SCROLL options have the same meanings as for a bound cursor.

An example:

```sql
OPEN curs1 FOR SELECT * FROM foo WHERE key = mykey;
```

### 43.7.2.2. `OPEN` FOR EXECUTE

```sql
OPEN unbound_cursorvar [ [ NO ] SCROLL ] FOR EXECUTE query_string
   [ USING expression [, ... ] ];
```

The cursor variable is opened and given the specified query to execute. The cursor cannot be open already, and it must have been declared as an unbound cursor variable (that is, as a simple refcursor variable). The query is specified as a string expression, in the same way as in the EXECUTE command. As usual, this gives flexibility so the query plan can vary from one run to the next (see Section 43.11.2), and it also means that variable substitution is not done on the command string. As with EXECUTE, parameter values can be inserted into the dynamic command via `format()` and `USING`. The SCROLL and NO SCROLL options have the same meanings as for a bound cursor.

An example:
OPEN curs1 FOR EXECUTE format('SELECT * FROM %I WHERE col1 = $1', tabname) USING keyvalue;

In this example, the table name is inserted into the query via format(). The comparison value for col1 is inserted via a USING parameter, so it needs no quoting.

43.7.2.3. Opening a Bound Cursor

OPEN bound_cursorvar [ ( [ argument_name := ] argument_value [, ...] ) ];

This form of OPEN is used to open a cursor variable whose query was bound to it when it was declared. The cursor cannot be open already. A list of actual argument value expressions must appear if and only if the cursor was declared to take arguments. These values will be substituted in the query.

The query plan for a bound cursor is always considered cacheable; there is no equivalent of EXECUTE in this case. Notice that SCROLL and NO SCROLL cannot be specified in OPEN, as the cursor's scrolling behavior was already determined.

Argument values can be passed using either positional or named notation. In positional notation, all arguments are specified in order. In named notation, each argument's name is specified using := to separate it from the argument expression. Similar to calling functions, described in Section 4.3, it is also allowed to mix positional and named notation.

Examples (these use the cursor declaration examples above):

OPEN curs2;
OPEN curs3(42);
OPEN curs3(key := 42);

Because variable substitution is done on a bound cursor's query, there are really two ways to pass values into the cursor: either with an explicit argument to OPEN, or implicitly by referencing a PL/pgSQL variable in the query. However, only variables declared before the bound cursor was declared will be substituted into it. In either case the value to be passed is determined at the time of the OPEN.

For example, another way to get the same effect as the curs3 example above is

DECLARE
  key integer;
BEGIN
  key := 42;
  OPEN curs4;

43.7.3. Using Cursors

Once a cursor has been opened, it can be manipulated with the statements described here.

These manipulations need not occur in the same function that opened the cursor to begin with. You can return a refcursor value out of a function and let the caller operate on the cursor. (Internally, a refcursor value is simply the string name of a so-called portal containing the active query for the cursor. This name can be passed around, assigned to other refcursor variables, and so on, without disturbing the portal.)

All portals are implicitly closed at transaction end. Therefore a refcursor value is usable to reference an open cursor only until the end of the transaction.
43.7.3.1. FETCH

FETCH [ direction { FROM | IN } ] cursor INTO target;

FETCH retrieves the next row from the cursor into a target, which might be a row variable, a record variable, or a comma-separated list of simple variables, just like SELECT INTO. If there is no next row, the target is set to NULL(s). As with SELECT INTO, the special variable FOUND can be checked to see whether a row was obtained or not.

The direction clause can be any of the variants allowed in the SQL FETCH command except the ones that can fetch more than one row; namely, it can be NEXT, PRIOR, FIRST, LAST, ABSOLUTE count, RELATIVE count, FORWARD, or BACKWARD. Omitting direction is the same as specifying NEXT. In the forms using a count, the count can be any integer-valued expression (unlike the SQL FETCH command, which only allows an integer constant). direction values that require moving backward are likely to fail unless the cursor was declared or opened with the SCROLL option.

cursor must be the name of a refcursor variable that references an open cursor portal.

Examples:

FETCH curs1 INTO rowvar;
FETCH curs2 INTO foo, bar, baz;
FETCH LAST FROM curs3 INTO x, y;
FETCH RELATIVE -2 FROM curs4 INTO x;

43.7.3.2. MOVE

MOVE [ direction { FROM | IN } ] cursor;

MOVE repositions a cursor without retrieving any data. MOVE works exactly like the FETCH command, except it only repositions the cursor and does not return the row moved to. As with SELECT INTO, the special variable FOUND can be checked to see whether there was a next row to move to.

Examples:

MOVE curs1;
MOVE LAST FROM curs3;
MOVE RELATIVE -2 FROM curs4;
MOVE FORWARD 2 FROM curs4;

43.7.3.3. UPDATE/DELETE WHERE CURRENT OF

UPDATE table SET ... WHERE CURRENT OF cursor;
DELETE FROM table WHERE CURRENT OF cursor;

When a cursor is positioned on a table row, that row can be updated or deleted using the cursor to identify the row. There are restrictions on what the cursor's query can be (in particular, no grouping) and it's best to use FOR UPDATE in the cursor. For more information see the DECLARE reference page.

An example:

UPDATE foo SET dataval = myval WHERE CURRENT OF curs1;
43.7.3.4. CLOSE

CLOSE cursor;

CLOSE closes the portal underlying an open cursor. This can be used to release resources earlier than end of transaction, or to free up the cursor variable to be opened again.

An example:

CLOSE curs1;

43.7.3.5. Returning Cursors

PL/pgSQL functions can return cursors to the caller. This is useful to return multiple rows or columns, especially with very large result sets. To do this, the function opens the cursor and returns the cursor name to the caller (or simply opens the cursor using a portal name specified by or otherwise known to the caller). The caller can then fetch rows from the cursor. The cursor can be closed by the caller, or it will be closed automatically when the transaction closes.

The portal name used for a cursor can be specified by the programmer or automatically generated. To specify a portal name, simply assign a string to the refcursor variable before opening it. The string value of the refcursor variable will be used by OPEN as the name of the underlying portal. However, if the refcursor variable is null, OPEN automatically generates a name that does not conflict with any existing portal, and assigns it to the refcursor variable.

Note

A bound cursor variable is initialized to the string value representing its name, so that the portal name is the same as the cursor variable name, unless the programmer overrides it by assignment before opening the cursor. But an unbound cursor variable defaults to the null value initially, so it will receive an automatically-generated unique name, unless overridden.

The following example shows one way a cursor name can be supplied by the caller:

CREATE TABLE test (col text);
INSERT INTO test VALUES ('123');
CREATE FUNCTION reffunc(refcursor) RETURNS refcursor AS '
BEGIN
OPEN $1 FOR SELECT col FROM test;
RETURN $1;
END;
' LANGUAGE plpgsql;
BEGIN;
SELECT reffunc('funccursor');
FETCH ALL IN funccursor;
COMMIT;

The following example uses automatic cursor name generation:

CREATE FUNCTION reffunc2() RETURNS refcursor AS '
DECLARE
ref refcursor;
BEGIN
  OPEN ref FOR SELECT col FROM test;
  RETURN ref;
END;
' LANGUAGE plpgsql;

-- need to be in a transaction to use cursors.
BEGIN;
SELECT reffunc2();

reffunc2
-------------------
<unnamed cursor 1>
(1 row)
FETCH ALL IN "<unnamed cursor 1>";
COMMIT;

The following example shows one way to return multiple cursors from a single function:

CREATE FUNCTION myfunc(refcursor, refcursor) RETURNS SETOF refcursor AS $$
BEGIN
  OPEN $1 FOR SELECT * FROM table_1;
  RETURN NEXT $1;
  OPEN $2 FOR SELECT * FROM table_2;
  RETURN NEXT $2;
END;
$$ LANGUAGE plpgsql;

-- need to be in a transaction to use cursors.
BEGIN;
SELECT * FROM myfunc('a', 'b');
FETCH ALL FROM a;
FETCH ALL FROM b;
COMMIT;

43.7.4. Looping Through a Cursor's Result

There is a variant of the FOR statement that allows iterating through the rows returned by a cursor. The syntax is:

[ <<label>> ]
FOR recordvar IN bound_cursorvar [ [ [ argument_name
  statements
END LOOP [ label ];

The cursor variable must have been bound to some query when it was declared, and it cannot be open already. The FOR statement automatically opens the cursor, and it closes the cursor again when the loop exits. A list of actual argument value expressions must appear if and only if the cursor was declared to take arguments. These values will be substituted in the query, in just the same way as during an OPEN (see Section 43.7.2.3).
The variable `recordvar` is automatically defined as type `record` and exists only inside the loop (any existing definition of the variable name is ignored within the loop). Each row returned by the cursor is successively assigned to this record variable and the loop body is executed.

### 43.8. Transaction Management

In procedures invoked by the `CALL` command as well as in anonymous code blocks (`DO` command), it is possible to end transactions using the commands `COMMIT` and `ROLLBACK`. A new transaction is started automatically after a transaction is ended using these commands, so there is no separate `START TRANSACTION` command. (Note that `BEGIN` and `END` have different meanings in PL/pgSQL.)

Here is a simple example:

```sql
CREATE PROCEDURE transaction_test1()
LANGUAGE plpgsql
AS $$
BEGIN
  FOR i IN 0..9 LOOP
    INSERT INTO test1 (a) VALUES (i);
    IF i % 2 = 0 THEN
      COMMIT;
    ELSE
      ROLLBACK;
    END IF;
  END LOOP;
END $$;
CALL transaction_test1();
```

Transaction control is only possible in `CALL` or `DO` invocations from the top level or nested `CALL` or `DO` invocations without any other intervening command. For example, if the call stack is `CALL proc1() -> CALL proc2() -> CALL proc3()`, then the second and third procedures can perform transaction control actions. But if the call stack is `CALL proc1() -> SELECT func2() -> CALL proc3()`, then the last procedure cannot do transaction control, because of the `SELECT` in between.

Special considerations apply to cursor loops. Consider this example:

```sql
CREATE PROCEDURE transaction_test2()
LANGUAGE plpgsql
AS $$
DECLARE
  r RECORD;
BEGIN
  FOR r IN SELECT * FROM test2 ORDER BY x LOOP
    INSERT INTO test1 (a) VALUES (r.x);
    COMMIT;
  END LOOP;
END $$;
CALL transaction_test2();
```

Normally, cursors are automatically closed at transaction commit. However, a cursor created as part of a loop like this is automatically converted to a holdable cursor by the first `COMMIT` or `ROLLBACK`.
That means that the cursor is fully evaluated at the first COMMIT or ROLLBACK rather than row by row. The cursor is still removed automatically after the loop, so this is mostly invisible to the user.

Transaction commands are not allowed in cursor loops driven by commands that are not read-only (for example UPDATE ... RETURNING).

A transaction cannot be ended inside a block with exception handlers.

## 43.9. Errors and Messages

### 43.9.1. Reporting Errors and Messages

Use the RAISE statement to report messages and raise errors.

```sql
RAISE [ level ] 'format' [, expression [, ... ]] [ USING option = expression [, ... ] ];
RAISE [ level ] condition_name [ USING option = expression [, ... ] ];
RAISE [ level ] SQLSTATE 'sqlstate' [ USING option = expression [, ... ] ];
RAISE [ level ] USING option = expression [, ... ];
RAISE ;
```

The `level` option specifies the error severity. Allowed levels are DEBUG, LOG, INFO, NOTICE, WARNING, and EXCEPTION, with EXCEPTION being the default. EXCEPTION raises an error (which normally aborts the current transaction); the other levels only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the log_min_messages and client_min_messages configuration variables. See Chapter 19 for more information.

After `level` if any, you can write a `format` (which must be a simple string literal, not an expression). The format string specifies the error message text to be reported. The format string can be followed by optional argument expressions to be inserted into the message. Inside the format string, `%` is replaced by the string representation of the next optional argument's value. Write `%%` to emit a literal `%`. The number of arguments must match the number of `%` placeholders in the format string, or an error is raised during the compilation of the function.

In this example, the value of `v_job_id` will replace the `%` in the string:

```sql
RAISE NOTICE 'Calling cs_create_job(%)', v_job_id;
```

You can attach additional information to the error report by writing `USING` followed by `option = expression` items. Each `expression` can be any string-valued expression. The allowed `option` key words are:

- **MESSAGE**
  Sets the error message text. This option can't be used in the form of RAISE that includes a format string before USING.

- **DETAIL**
  Supplies an error detail message.

- **HINT**
  Supplies a hint message.
ERRCODE

Specifies the error code (SQLSTATE) to report, either by condition name, as shown in Appendix A, or directly as a five-character SQLSTATE code.

COLUMN
CONSTRAINT
DATATYPE
TABLE
SCHEMA

Supplies the name of a related object.

This example will abort the transaction with the given error message and hint:

RAISE EXCEPTION 'Nonexistent ID --> %', user_id
  USING HINT = 'Please check your user ID';

These two examples show equivalent ways of setting the SQLSTATE:

RAISE 'Duplicate user ID: %', user_id USING ERRCODE = 'unique_violation';
RAISE 'Duplicate user ID: %', user_id USING ERRCODE = '23505';

There is a second RAISE syntax in which the main argument is the condition name or SQLSTATE to be reported, for example:

RAISE division_by_zero;
RAISE SQLSTATE '22012';

In this syntax, USING can be used to supply a custom error message, detail, or hint. Another way to do the earlier example is

RAISE unique_violation USING MESSAGE = 'Duplicate user ID: ' || user_id;

Still another variant is to write RAISE USING or RAISE level USING and put everything else into the USING list.

The last variant of RAISE has no parameters at all. This form can only be used inside a BEGIN block's EXCEPTION clause; it causes the error currently being handled to be re-thrown.

Note

Before PostgreSQL 9.1, RAISE without parameters was interpreted as re-throwing the error from the block containing the active exception handler. Thus an EXCEPTION clause nested within that handler could not catch it, even if the RAISE was within the nested EXCEPTION clause's block. This was deemed surprising as well as being incompatible with Oracle's PL/SQL.

If no condition name nor SQLSTATE is specified in a RAISE EXCEPTION command, the default is to use RAISE_EXCEPTION (P0001). If no message text is specified, the default is to use the condition name or SQLSTATE as message text.
Note

When specifying an error code by SQLSTATE code, you are not limited to the pre-defined error codes, but can select any error code consisting of five digits and/or upper-case ASCII letters, other than 00000. It is recommended that you avoid throwing error codes that end in three zeroes, because these are category codes and can only be trapped by trapping the whole category.

43.9.2. Checking Assertions

The `ASSERT` statement is a convenient shorthand for inserting debugging checks into PL/pgSQL functions.

```
ASSERT condition [ , message ];
```

The `condition` is a Boolean expression that is expected to always evaluate to true; if it does, the `ASSERT` statement does nothing further. If the result is false or null, then an `ASSERT_FAILURE` exception is raised. (If an error occurs while evaluating the `condition`, it is reported as a normal error.)

If the optional `message` is provided, it is an expression whose result (if not null) replaces the default error message text “assertion failed”, should the `condition` fail. The `message` expression is not evaluated in the normal case where the assertion succeeds.

Testing of assertions can be enabled or disabled via the configuration parameter `plpgsql.check_asserts`, which takes a Boolean value; the default is on. If this parameter is off then `ASSERT` statements do nothing.

Note that `ASSERT` is meant for detecting program bugs, not for reporting ordinary error conditions. Use the `RAISE` statement, described above, for that.

43.10. Trigger Functions

PL/pgSQL can be used to define trigger functions on data changes or database events. A trigger function is created with the `CREATE FUNCTION` command, declaring it as a function with no arguments and a return type of `trigger` (for data change triggers) or `event_trigger` (for database event triggers). Special local variables named `TG_something` are automatically defined to describe the condition that triggered the call.

43.10.1. Triggers on Data Changes

A data change trigger is declared as a function with no arguments and a return type of `trigger`. Note that the function must be declared with no arguments even if it expects to receive some arguments specified in `CREATE TRIGGER` — such arguments are passed via `TG_ARGV`, as described below.

When a PL/pgSQL function is called as a trigger, several special variables are created automatically in the top-level block. They are:

```
NEW
```

Data type `RECORD`; variable holding the new database row for `INSERT/UPDATE` operations in row-level triggers. This variable is unassigned in statement-level triggers and for `DELETE` operations.
OLD

Data type RECORD; variable holding the old database row for UPDATE/DELETE operations in row-level triggers. This variable is unassigned in statement-level triggers and for INSERT operations.

TG_NAME

Data type name; variable that contains the name of the trigger actually fired.

TG_WHEN

Data type text; a string of BEFORE, AFTER, or INSTEAD OF, depending on the trigger’s definition.

TG_LEVEL

Data type text; a string of either ROW or STATEMENT depending on the trigger’s definition.

TG_OP

Data type text; a string of INSERT, UPDATE, DELETE, or TRUNCATE telling for which operation the trigger was fired.

TG_RELID

Data type oid; the object ID of the table that caused the trigger invocation.

TG_RELNAME

Data type name; the name of the table that caused the trigger invocation. This is now deprecated, and could disappear in a future release. Use TG_TABLE_NAME instead.

TG_TABLE_NAME

Data type name; the name of the table that caused the trigger invocation.

TG_TABLE_SCHEMA

Data type name; the name of the schema of the table that caused the trigger invocation.

TG_NARGS

Data type integer; the number of arguments given to the trigger function in the CREATE TRIGGER statement.

TG_ARGV[]

Data type array of text; the arguments from the CREATE TRIGGER statement. The index counts from 0. Invalid indexes (less than 0 or greater than or equal to tg_nargs) result in a null value.

A trigger function must return either NULL or a record/row value having exactly the structure of the table the trigger was fired for.

Row-level triggers fired BEFORE can return null to signal the trigger manager to skip the rest of the operation for this row (i.e., subsequent triggers are not fired, and the INSERT/UPDATE/DELETE does not occur for this row). If a nonnull value is returned then the operation proceeds with that row value. Returning a row value different from the original value of NEW alters the row that will be inserted or updated. Thus, if the trigger function wants the triggering action to succeed normally without altering...
the row value, NEW (or a value equal thereto) has to be returned. To alter the row to be stored, it is possible to replace single values directly in NEW and return the modified NEW, or to build a complete new record/row to return. In the case of a before-trigger on DELETE, the returned value has no direct effect, but it has to be nonnull to allow the trigger action to proceed. Note that NEW is null in DELETE triggers, so returning that is usually not sensible. The usual idiom in DELETE triggers is to return OLD.

INSTEAD OF triggers (which are always row-level triggers, and may only be used on views) can return null to signal that they did not perform any updates, and that the rest of the operation for this row should be skipped (i.e., subsequent triggers are not fired, and the row is not counted in the row-affected status for the surrounding INSERT/UPDATE/DELETE). Otherwise a nonnull value should be returned, to signal that the trigger performed the requested operation. For INSERT and UPDATE operations, the return value should be NEW, which the trigger function may modify to support INSERT RETURNING and UPDATE RETURNING (this will also affect the row value passed to any subsequent triggers, or passed to a special EXCLUDED alias reference within an INSERT statement with an ON CONFLICT DO UPDATE clause). For DELETE operations, the return value should be OLD.

The return value of a row-level trigger fired AFTER or a statement-level trigger fired BEFORE or AFTER is always ignored; it might as well be null. However, any of these types of triggers might still abort the entire operation by raising an error.

Example 43.3 shows an example of a trigger function in PL/pgSQL.

Example 43.3. A PL/pgSQL Trigger Function

This example trigger ensures that any time a row is inserted or updated in the table, the current user name and time are stamped into the row. And it checks that an employee's name is given and that the salary is a positive value.

CREATE TABLE emp (
    empname text,
    salary integer,
    last_date timestamp,
    last_user text
);

CREATE FUNCTION emp_stamp() RETURNS trigger AS $emp_stamp$
BEGIN
    -- Check that empname and salary are given
    IF NEW.empname IS NULL THEN
        RAISE EXCEPTION 'empname cannot be null';
    END IF;
    IF NEW.salary IS NULL THEN
        RAISE EXCEPTION '% cannot have null salary',
    END IF;
    NEW.empname;
END IF;

    -- Who works for us when they must pay for it?
    IF NEW.salary < 0 THEN
        RAISE EXCEPTION '% cannot have a negative salary',
    END IF;
    NEW.empname;
END IF;

    -- Remember who changed the payroll when
    NEW.last_date := current_timestamp;
    NEW.last_user := current_user;
    RETURN NEW;
END;
$emp_stamp$ LANGUAGE plpgsql;
CREATE TRIGGER emp_stamp BEFORE INSERT OR UPDATE ON emp
    FOR EACH ROW EXECUTE FUNCTION emp_stamp();

Another way to log changes to a table involves creating a new table that holds a row for each insert, update, or delete that occurs. This approach can be thought of as auditing changes to a table. Example 43.4 shows an example of an audit trigger function in PL/pgSQL.

**Example 43.4. A PL/pgSQL Trigger Function For Auditing**

This example trigger ensures that any insert, update or delete of a row in the `emp` table is recorded (i.e., audited) in the `emp_audit` table. The current time and user name are stamped into the row, together with the type of operation performed on it.

```sql
CREATE TABLE emp (
    empname           text NOT NULL,
    salary            integer
);
CREATE TABLE emp_audit(
    operation         char(1)   NOT NULL,
    stamp             timestamp NOT NULL,
    userid            text      NOT NULL,
    empname           text      NOT NULL,
    salary integer
);
CREATE OR REPLACE FUNCTION process_emp_audit() RETURNS TRIGGER AS
$emp_audit$
BEGIN
    -- Create a row in emp_audit to reflect the operation
    -- performed on emp,
    -- making use of the special variable TG_OP to work out the
    -- operation.
    --
    IF (TG_OP = 'DELETE') THEN
        INSERT INTO emp_audit SELECT 'D', now(), user, OLD.*;
    ELSIF (TG_OP = 'UPDATE') THEN
        INSERT INTO emp_audit SELECT 'U', now(), user, NEW.*;
    ELSIF (TG_OP = 'INSERT') THEN
        INSERT INTO emp_audit SELECT 'I', now(), user, NEW.*;
    END IF;
    RETURN NULL; -- result is ignored since this is an AFTER
    trigger
END;
$emp_audit$ LANGUAGE plpgsql;
CREATE TRIGGER emp_audit
AFTER INSERT OR UPDATE OR DELETE ON emp
    FOR EACH ROW EXECUTE FUNCTION process_emp_audit();
```

A variation of the previous example uses a view joining the main table to the audit table, to show when each entry was last modified. This approach still records the full audit trail of changes to the table, but also presents a simplified view of the audit trail, showing just the last modified timestamp derived from the audit trail for each entry. Example 43.5 shows an example of an audit trigger on a view in PL/pgSQL.
Example 43.5. A PL/pgSQL View Trigger Function For Auditing

This example uses a trigger on the view to make it updatable, and ensure that any insert, update or delete of a row in the view is recorded (i.e., audited) in the emp_audit table. The current time and user name are recorded, together with the type of operation performed, and the view displays the last modified time of each row.

CREATE TABLE emp (  
    empname text PRIMARY KEY,  
    salary integer  
);  
CREATE TABLE emp_audit (  
    operation char(1) NOT NULL,  
    userid text NOT NULL,  
    empname text NOT NULL,  
    salary integer,  
    stamp timestamp NOT NULL  
);  
CREATE VIEW emp_view AS  
    SELECT e.empname,  
           e.salary,  
           max(ea.stamp) AS last_updated  
    FROM emp e  
    LEFT JOIN emp_audit ea ON ea.empname = e.empname  
    GROUP BY 1, 2;  
CREATE OR REPLACE FUNCTION update_emp_view() RETURNS TRIGGER AS $$  
BEGIN  
    -- Perform the required operation on emp, and create a row  
    -- in emp_audit  
    -- to reflect the change made to emp.  
    --  
    IF (TG_OP = 'DELETE') THEN  
        DELETE FROM emp WHERE empname = OLD.empname;  
        IF NOT FOUND THEN RETURN NULL; END IF;  
        OLD.last_updated = now();  
        INSERT INTO emp_audit VALUES('D', user, OLD.*);  
        RETURN OLD;  
    ELSIF (TG_OP = 'UPDATE') THEN  
        UPDATE emp SET salary = NEW.salary WHERE empname =  
        OLD.empname;  
        IF NOT FOUND THEN RETURN NULL; END IF;  
        NEW.last_updated = now();  
        INSERT INTO emp_audit VALUES('U', user, NEW.*);  
        RETURN NEW;  
    ELSIF (TG_OP = 'INSERT') THEN  
        INSERT INTO emp VALUES(NEW.empname, NEW.salary);  
        NEW.last_updated = now();  
        INSERT INTO emp_audit VALUES('I', user, NEW.*);  
        RETURN NEW;  
    END IF;  
$$;
One use of triggers is to maintain a summary table of another table. The resulting summary can be used in place of the original table for certain queries — often with vastly reduced run times. This technique is commonly used in Data Warehousing, where the tables of measured or observed data (called fact tables) might be extremely large. Example 43.6 shows an example of a trigger function in PL/pgSQL that maintains a summary table for a fact table in a data warehouse.

Example 43.6. A PL/pgSQL Trigger Function For Maintaining A Summary Table

The schema detailed here is partly based on the *Grocery Store* example from *The Data Warehouse Toolkit* by Ralph Kimball.

```
CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
    FOR EACH ROW EXECUTE FUNCTION update_emp_view();
```

```sql
PL/pgSQL - SQL Pro-
cedural Language
END;
$$ LANGUAGE plpgsql;
```

```
CREATE TABLE time_dimension (  
    time_key integer NOT NULL,
    day_of_week integer NOT NULL,
    day_of_month integer NOT NULL,
    month integer NOT NULL,
    quarter integer NOT NULL,
    year integer NOT NULL
);
CREATE UNIQUE INDEX time_dimension_key ON time_dimension(time_key);
CREATE TABLE sales_fact (  
    time_key integer NOT NULL,
    product_key integer NOT NULL,
    store_key integer NOT NULL,
    amount_sold numeric(12,2) NOT NULL,
    units_sold integer NOT NULL,
    amount_cost numeric(12,2) NOT NULL
);
CREATE INDEX sales_fact_time ON sales_fact(time_key);
```

```
CREATE TABLE sales_summary_bytime (  
    time_key integer NOT NULL,
    amount_sold numeric(15,2) NOT NULL,
    units_sold numeric(12) NOT NULL,
    amount_cost numeric(15,2) NOT NULL
);
CREATE UNIQUE INDEX sales_summary_bytime_key ON
    sales_summary_bytime(time_key);
```

```
CREATE TABLE sales_summary_bytime (  
    time_key integer NOT NULL,
    amount_sold numeric(15,2) NOT NULL,
    units_sold numeric(12) NOT NULL,
    amount_cost numeric(15,2) NOT NULL
);
CREATE UNIQUE INDEX sales_summary_bytime_key ON
    sales_summary_bytime(time_key);
```

```
CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
    FOR EACH ROW EXECUTE FUNCTION update_emp_view();
```

```
CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
    FOR EACH ROW EXECUTE FUNCTION update_emp_view();
```

Example 43.6. A PL/pgSQL Trigger Function For Maintaining A Summary Table

The schema detailed here is partly based on the *Grocery Store* example from *The Data Warehouse Toolkit* by Ralph Kimball.

```
CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
    FOR EACH ROW EXECUTE FUNCTION update_emp_view();
```

```
CREATE TABLE time_dimension (  
    time_key integer NOT NULL,
    day_of_week integer NOT NULL,
    day_of_month integer NOT NULL,
    month integer NOT NULL,
    quarter integer NOT NULL,
    year integer NOT NULL
);
CREATE UNIQUE INDEX time_dimension_key ON time_dimension(time_key);
CREATE TABLE sales_fact (  
    time_key integer NOT NULL,
    product_key integer NOT NULL,
    store_key integer NOT NULL,
    amount_sold numeric(12,2) NOT NULL,
    units_sold integer NOT NULL,
    amount_cost numeric(12,2) NOT NULL
);
CREATE INDEX sales_fact_time ON sales_fact(time_key);
```

```
CREATE TABLE sales_summary_bytime (  
    time_key integer NOT NULL,
    amount_sold numeric(15,2) NOT NULL,
    units_sold numeric(12) NOT NULL,
    amount_cost numeric(15,2) NOT NULL
);
CREATE UNIQUE INDEX sales_summary_bytime_key ON
    sales_summary_bytime(time_key);
```

```
CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
    FOR EACH ROW EXECUTE FUNCTION update_emp_view();
```
CREATE OR REPLACE FUNCTION maint_sales_summary_bytime() RETURNS TRIGGER AS $maint_sales_summary_bytime$
DECLARE
  delta_time_key          integer;
  delta_amount_sold       numeric(15,2);
  delta_units_sold        numeric(12);
  delta_amount_cost       numeric(15,2);
BEGIN

  -- Work out the increment/decrement amount(s).
  IF (TG_OP = 'DELETE') THEN
    delta_time_key = OLD.time_key;
    delta_amount_sold = -1 * OLD.amount_sold;
    delta_units_sold = -1 * OLD.units_sold;
    delta_amount_cost = -1 * OLD.amount_cost;
  ELSIF (TG_OP = 'UPDATE') THEN
    -- forbid updates that change the time_key -
    -- (probably not too onerous, as DELETE + INSERT is how most
    -- changes will be made).
    IF ( OLD.time_key != NEW.time_key) THEN
      RAISE EXCEPTION 'Update of time_key : % -> % not allowed',
                         OLD.time_key, NEW.time_key;
    END IF;
    delta_time_key = OLD.time_key;
    delta_amount_sold = NEW.amount_sold - OLD.amount_sold;
    delta_units_sold = NEW.units_sold - OLD.units_sold;
    delta_amount_cost = NEW.amount_cost - OLD.amount_cost;
  ELSIF (TG_OP = 'INSERT') THEN
    delta_time_key = NEW.time_key;
    delta_amount_sold = NEW.amount_sold;
    delta_units_sold = NEW.units_sold;
    delta_amount_cost = NEW.amount_cost;
  END IF;

  -- Insert or update the summary row with the new values.
  LOOP
    UPDATE sales_summary_bytime
    SET amount_sold = amount_sold + delta_amount_sold,
        units_sold = units_sold + delta_units_sold,
        amount_cost = amount_cost + delta_amount_cost
    WHERE time_key = delta_time_key;
    EXIT insert_update WHEN found;
END;
BEGIN
    INSERT INTO sales_summary_bytime (time_key, amount_sold, units_sold, amount_cost)
    VALUES (delta_time_key, delta_amount_sold, delta_units_sold, delta_amount_cost);
    EXIT insert_update;

EXCEPTION
    WHEN UNIQUE_VIOLATION THEN
        -- do nothing
    END;
END LOOP insert_update;

RETURN NULL;
END;
$maint_sales_summary_bytime$ LANGUAGE plpgsql;

CREATE TRIGGER maint_sales_summary_bytime
AFTER INSERT OR UPDATE OR DELETE ON sales_fact
FOR EACH ROW EXECUTE FUNCTION maint_sales_summary_bytime();

INSERT INTO sales_fact VALUES(1,1,1,10,3,15);
INSERT INTO sales_fact VALUES(1,2,1,20,5,35);
INSERT INTO sales_fact VALUES(2,2,1,40,15,135);
INSERT INTO sales_fact VALUES(2,3,1,10,1,13);
SELECT * FROM sales_summary_bytime;
DELETE FROM sales_fact WHERE product_key = 1;
SELECT * FROM sales_summary_bytime;
UPDATE sales_fact SET units_sold = units_sold * 2;
SELECT * FROM sales_summary_bytime;

AFTER triggers can also make use of transition tables to inspect the entire set of rows changed by the triggering statement. The CREATE TRIGGER command assigns names to one or both transition tables, and then the function can refer to those names as though they were read-only temporary tables. Example 43.7 shows an example.

Example 43.7. Auditing with Transition Tables

This example produces the same results as Example 43.4, but instead of using a trigger that fires for every row, it uses a trigger that fires once per statement, after collecting the relevant information in a transition table. This can be significantly faster than the row-trigger approach when the invoking statement has modified many rows. Notice that we must make a separate trigger declaration for each kind of event, since the REFERENCING clauses must be different for each case. But this does not stop us from using a single trigger function if we choose. (In practice, it might be better to use three separate functions and avoid the run-time tests on TG_OP.)

CREATE TABLE emp (empname text NOT NULL, salary integer
CREATE TABLE emp_audit(
    operation         char(1)   NOT NULL,
    stamp             timestamp NOT NULL,
    userid            text      NOT NULL,
    empname           text      NOT NULL,
    salary integer
);

CREATE OR REPLACE FUNCTION process_emp_audit() RETURNS TRIGGER AS
$emp_audit$
BEGIN
    -- Create rows in emp_audit to reflect the operations performed on emp,
    -- making use of the special variable TG_OP to work out the operation.
    --
    -- IF (TG_OP = 'DELETE') THEN
    INSERT INTO emp_audit
    SELECT 'D', now(), user, o.* FROM old_table o;
    ELSIF (TG_OP = 'UPDATE') THEN
    INSERT INTO emp_audit
    SELECT 'U', now(), user, n.* FROM new_table n;
    ELSIF (TG_OP = 'INSERT') THEN
    INSERT INTO emp_audit
    SELECT 'I', now(), user, n.* FROM new_table n;
    END IF;
    RETURN NULL; -- result is ignored since this is an AFTER trigger
END;
$emp_audit$ LANGUAGE plpgsql;

CREATE TRIGGER emp_audit_ins
AFTER INSERT ON emp
REFERENCING NEW TABLE AS new_table
FOR EACH STATEMENT EXECUTE FUNCTION process_emp_audit();

CREATE TRIGGER emp_audit_upd
AFTER UPDATE ON emp
REFERENCING OLD TABLE AS old_table NEW TABLE AS new_table
FOR EACH STATEMENT EXECUTE FUNCTION process_emp_audit();

CREATE TRIGGER emp_audit_del
AFTER DELETE ON emp
REFERENCING OLD TABLE AS old_table
FOR EACH STATEMENT EXECUTE FUNCTION process_emp_audit();

43.10.2. Triggers on Events

PL/pgSQL can be used to define event triggers. PostgreSQL requires that a function that is to be called as an event trigger must be declared as a function with no arguments and a return type of event_trigger.

When a PL/pgSQL function is called as an event trigger, several special variables are created automatically in the top-level block. They are:

TG_EVENT

Data type text: a string representing the event the trigger is fired for.
TG_TAG

Data type text; variable that contains the command tag for which the trigger is fired.

Example 43.8 shows an example of an event trigger function in PL/pgSQL.

**Example 43.8. A PL/pgSQL Event Trigger Function**

This example trigger simply raises a NOTICE message each time a supported command is executed.

```
CREATE OR REPLACE FUNCTION snitch() RETURNS event_trigger AS $$
BEGIN
    RAISE NOTICE 'snitch: % %', tg_event, tg_tag;
END;
$$ LANGUAGE plpgsql;

CREATE EVENT TRIGGER snitch ON ddl_command_start EXECUTE FUNCTION
snitch();
```

**43.11. PL/pgSQL Under the Hood**

This section discusses some implementation details that are frequently important for PL/pgSQL users to know.

### 43.11.1. Variable Substitution

SQL statements and expressions within a PL/pgSQL function can refer to variables and parameters of the function. Behind the scenes, PL/pgSQL substitutes query parameters for such references. Parameters will only be substituted in places where a parameter or column reference is syntactically allowed. As an extreme case, consider this example of poor programming style:

```
INSERT INTO foo (foo) VALUES (foo);
```

The first occurrence of `foo` must syntactically be a table name, so it will not be substituted, even if the function has a variable named `foo`. The second occurrence must be the name of a column of the table, so it will not be substituted either. Only the third occurrence is a candidate to be a reference to the function’s variable.

**Note**

PostgreSQL versions before 9.0 would try to substitute the variable in all three cases, leading to syntax errors.

Since the names of variables are syntactically no different from the names of table columns, there can be ambiguity in statements that also refer to tables: is a given name meant to refer to a table column, or a variable? Let's change the previous example to

```
INSERT INTO dest (col) SELECT foo + bar FROM src;
```

Here, `dest` and `src` must be table names, and `col` must be a column of `dest`, but `foo` and `bar` might reasonably be either variables of the function or columns of `src`.

By default, PL/pgSQL will report an error if a name in a SQL statement could refer to either a variable or a table column. You can fix such a problem by renaming the variable or column, or by qualifying the ambiguous reference, or by telling PL/pgSQL which interpretation to prefer.
The simplest solution is to rename the variable or column. A common coding rule is to use a different naming convention for PL/pgSQL variables than you use for column names. For example, if you consistently name function variables `v_something` while none of your column names start with `v_`, no conflicts will occur.

Alternatively you can qualify ambiguous references to make them clear. In the above example, `src.foo` would be an unambiguous reference to the table column. To create an unambiguous reference to a variable, declare it in a labeled block and use the block’s label (see Section 43.2). For example,

```
<<block>>
DECLARE
  foo int;
BEGIN
  foo := ...;
  INSERT INTO dest (col) SELECT block.foo + bar FROM src;
```

Here `block.foo` means the variable even if there is a column `foo` in `src`. Function parameters, as well as special variables such as `FOUND`, can be qualified by the function’s name, because they are implicitly declared in an outer block labeled with the function’s name.

Sometimes it is impractical to fix all the ambiguous references in a large body of PL/pgSQL code. In such cases you can specify that PL/pgSQL should resolve ambiguous references as the variable (which is compatible with PL/pgSQL’s behavior before PostgreSQL 9.0), or as the table column (which is compatible with some other systems such as Oracle).

To change this behavior on a system-wide basis, set the configuration parameter `plpgsql.variable_conflict` to one of `error`, `use_variable`, or `use_column` (where `error` is the factory default). This parameter affects subsequent compilations of statements in PL/pgSQL functions, but not statements already compiled in the current session. Because changing this setting can cause unexpected changes in the behavior of PL/pgSQL functions, it can only be changed by a superuser.

You can also set the behavior on a function-by-function basis, by inserting one of these special commands at the start of the function text:

```
#variable_conflict error
#variable_conflict use_variable
#variable_conflict use_column
```

These commands affect only the function they are written in, and override the setting of `plpgsql.variable_conflict`. An example is

```
CREATE FUNCTION stamp_user(id int, comment text) RETURNS void AS $$
  #variable_conflict use_variable
  DECLARE
    curtime timestamp := now();
  BEGIN
    UPDATE users SET last_modified = curtime, comment = comment
      WHERE users.id = id;
  END;
$$ LANGUAGE plpgsql;
```

In the `UPDATE` command, `curtime`, `comment`, and `id` will refer to the function’s variable and parameters whether or not `users` has columns of those names. Notice that we had to qualify the reference to `users.id` in the `WHERE` clause to make it refer to the table column. But we did not have to qualify the reference to `comment` as a target in the `UPDATE` list, because syntactically that must be
a column of users. We could write the same function without depending on the variable_con-

flict setting in this way:

```
CREATE FUNCTION stamp_user(id int, comment text) RETURNS void AS $$
<<fn>>
DECLARE
curtime timestamp := now();
BEGIN
UPDATE users SET last_modified = fn.curtime, comment =
stamp_user.comment
WHERE users.id = stamp_user.id;
END;
$$ LANGUAGE plpgsql;
```

Variable substitution does not happen in the command string given to EXECUTE or one of its variants. If you need to insert a varying value into such a command, do so as part of constructing the string value, or use USING, as illustrated in Section 43.5.4.

Variable substitution currently works only in SELECT, INSERT, UPDATE, and DELETE commands, because the main SQL engine allows query parameters only in these commands. To use a non-constant name or value in other statement types (generically called utility statements), you must construct the utility statement as a string and EXECUTE it.

### 43.11.2. Plan Caching

The PL/pgSQL interpreter parses the function’s source text and produces an internal binary instruction tree the first time the function is called (within each session). The instruction tree fully translates the PL/pgSQL statement structure, but individual SQL expressions and SQL commands used in the function are not translated immediately.

As each expression and SQL command is first executed in the function, the PL/pgSQL interpreter parses and analyzes the command to create a prepared statement, using the SPI manager's SPI_prepare function. Subsequent visits to that expression or command reuse the prepared statement. Thus, a function with conditional code paths that are seldom visited will never incur the overhead of analyzing those commands that are never executed within the current session. A disadvantage is that errors in a specific expression or command cannot be detected until that part of the function is reached in execution. (Trivial syntax errors will be detected during the initial parsing pass, but anything deeper will not be detected until execution.)

PL/pgSQL (or more precisely, the SPI manager) can furthermore attempt to cache the execution plan associated with any particular prepared statement. If a cached plan is not used, then a fresh execution plan is generated on each visit to the statement, and the current parameter values (that is, PL/pgSQL variable values) can be used to optimize the selected plan. If the statement has no parameters, or is executed many times, the SPI manager will consider creating a generic plan that is not dependent on specific parameter values, and caching that for re-use. Typically this will happen only if the execution plan is not very sensitive to the values of the PL/pgSQL variables referenced in it. If it is, generating a plan each time is a net win. See PREPARE for more information about the behavior of prepared statements.

Because PL/pgSQL saves prepared statements and sometimes execution plans in this way, SQL commands that appear directly in a PL/pgSQL function must refer to the same tables and columns on every execution; that is, you cannot use a parameter as the name of a table or column in an SQL command. To get around this restriction, you can construct dynamic commands using the PL/pgSQL EXECUTE statement — at the price of performing new parse analysis and constructing a new execution plan on every execution.

The mutable nature of record variables presents another problem in this connection. When fields of a record variable are used in expressions or statements, the data types of the fields must not change.
from one call of the function to the next, since each expression will be analyzed using the data type that is present when the expression is first reached. EXECUTE can be used to get around this problem when necessary.

If the same function is used as a trigger for more than one table, PL/pgSQL prepares and caches statements independently for each such table — that is, there is a cache for each trigger function and table combination, not just for each function. This alleviates some of the problems with varying data types; for instance, a trigger function will be able to work successfully with a column named *key* even if it happens to have different types in different tables.

Likewise, functions having polymorphic argument types have a separate statement cache for each combination of actual argument types they have been invoked for, so that data type differences do not cause unexpected failures.

Statement caching can sometimes have surprising effects on the interpretation of time-sensitive values. For example there is a difference between what these two functions do:

```sql
CREATE FUNCTION logfunc1(logtxt text) RETURNS void AS $$
BEGIN
    INSERT INTO logtable VALUES (logtxt, 'now');
END;
$$ LANGUAGE plpgsql;
```

and:

```sql
CREATE FUNCTION logfunc2(logtxt text) RETURNS void AS $$
DECLARE
    curtime timestamp;
BEGIN
    curtime := 'now';
    INSERT INTO logtable VALUES (logtxt, curtime);
END;
$$ LANGUAGE plpgsql;
```

In the case of `logfunc1`, the PostgreSQL main parser knows when analyzing the `INSERT` that the string 'now' should be interpreted as `timestamp`, because the target column of `logtable` is of that type. Thus, 'now' will be converted to a `timestamp` constant when the `INSERT` is analyzed, and then used in all invocations of `logfunc1` during the lifetime of the session. Needless to say, this isn't what the programmer wanted. A better idea is to use the `now()` or `current_timestamp` function.

In the case of `logfunc2`, the PostgreSQL main parser does not know what type 'now' should become and therefore it returns a data value of type `text` containing the string `now`. During the ensuing assignment to the local variable `curtime`, the PL/pgSQL interpreter casts this string to the `timestamp` type by calling the `text_out` and `timestamp_in` functions for the conversion. So, the computed time stamp is updated on each execution as the programmer expects. Even though this happens to work as expected, it’s not terribly efficient, so use of the `now()` function would still be a better idea.

### 43.12. Tips for Developing in PL/pgSQL

One good way to develop in PL/pgSQL is to use the text editor of your choice to create your functions, and in another window, use psql to load and test those functions. If you are doing it this way, it is a good idea to write the function using `CREATE OR REPLACE FUNCTION`. That way you can just reload the file to update the function definition. For example:
CREATE OR REPLACE FUNCTION testfunc(integer) RETURNS integer AS $$
....
$$ LANGUAGE plpgsql;

While running psql, you can load or reload such a function definition file with:

\i filename.sql

and then immediately issue SQL commands to test the function.

Another good way to develop in PL/pgSQL is with a GUI database access tool that facilitates development in a procedural language. One example of such a tool is pgAdmin, although others exist. These tools often provide convenient features such as escaping single quotes and making it easier to recreate and debug functions.

43.12.1. Handling of Quotation Marks

The code of a PL/pgSQL function is specified in CREATE FUNCTION as a string literal. If you write the string literal in the ordinary way with surrounding single quotes, then any single quotes inside the function body must be doubled; likewise any backslashes must be doubled (assuming escape string syntax is used). Doubling quotes is at best tedious, and in more complicated cases the code can become downright incomprehensible, because you can easily find yourself needing half a dozen or more adjacent quote marks. It’s recommended that you instead write the function body as a “dollar-quoted” string literal (see Section 4.1.2.4). In the dollar-quoted approach, you never double any quote marks, but instead take care to choose a different dollar-quoted delimiter for each level of nesting you need. For example, you might write the CREATE FUNCTION command as:

CREATE OR REPLACE FUNCTION testfunc(integer) RETURNS integer AS
$PROC$
....
$PROC$ LANGUAGE plpgsql;

Within this, you might use quote marks for simple literal strings in SQL commands and $$ to delimit fragments of SQL commands that you are assembling as strings. If you need to quote text that includes $$, you could use $Q$, and so on.

The following chart shows what you have to do when writing quote marks without dollar quoting. It might be useful when translating pre-dollar quoting code into something more comprehensible.

1 quotation mark

To begin and end the function body, for example:

CREATE FUNCTION foo() RETURNS integer AS ' 
....
' LANGUAGE plpgsql;

Anywhere within a single-quoted function body, quote marks must appear in pairs.

2 quotation marks

For string literals inside the function body, for example:

a_output := ''Blah'';
SELECT * FROM users WHERE f_name=''foobar'';
In the dollar-quoting approach, you'd just write:

```sql
a_output := 'Blah';
SELECT * FROM users WHERE f_name='foobar';
```

which is exactly what the PL/pgSQL parser would see in either case.

**4 quotation marks**

When you need a single quotation mark in a string constant inside the function body, for example:

```sql
a_output := a_output || '' AND name LIKE '''foobar'''' AND xyz''
```

The value actually appended to `a_output` would be:

```
AND name LIKE 'foobar' AND xyz.
```

In the dollar-quoting approach, you'd write:

```sql
a_output := a_output || $$ AND name LIKE 'foobar' AND xyz$$
```

being careful that any dollar-quote delimiters around this are not just $$.  

**6 quotation marks**

When a single quotation mark in a string inside the function body is adjacent to the end of that string constant, for example:

```sql
a_output := a_output || '' AND name LIKE '''foobar''''''
```

The value appended to `a_output` would then be:

```
AND name LIKE 'foobar'.
```

In the dollar-quoting approach, this becomes:

```sql
a_output := a_output || $$ AND name LIKE 'foobar'$$
```

**10 quotation marks**

When you want two single quotation marks in a string constant (which accounts for 8 quotation marks) and this is adjacent to the end of that string constant (2 more). You will probably only need that if you are writing a function that generates other functions, as in Example 43.10. For example:

```sql
a_output := a_output || '' if v_'' ||
referer_keys.kind || '' like ''''''''''
|| referer_keys.key_string || ''''''''''
then return '''''''''' || referer_keys.referrer_type
|| ''''''; end if;'';
```

The value of `a_output` would then be:

```sql
if v_... like '''...''' then return '''...''' end if;
```

In the dollar-quoting approach, this becomes:
```
a_output := a_output || $$ if v_$$ || referrer_keys.kind || $$
     || referrer_keys.key_string || $$'
then return '$$ || referrer_keys.referrer_type
     || $$'; end if;$$;
```

where we assume we only need to put single quote marks into `a_output`, because it will be re-quoted before use.

### 43.12.2. Additional Compile-time Checks

To aid the user in finding instances of simple but common problems before they cause harm, PL/pgSQL provides additional checks. When enabled, depending on the configuration, they can be used to emit either a WARNING or an ERROR during the compilation of a function. A function which has received a WARNING can be executed without producing further messages, so you are advised to test in a separate development environment.

These additional checks are enabled through the configuration variables `plpgsql.extra_warnings` for warnings and `plpgsql.extra_errors` for errors. Both can be set either to a comma-separated list of checks, "none" or "all". The default is "none". Currently the list of available checks includes only one:

```
shadowed_variables
```

Checks if a declaration shadows a previously defined variable.

The following example shows the effect of `plpgsql.extra_warnings` set to `shadowed_variables`:

```
SET plpgsql.extra_warnings TO 'shadowed_variables';

CREATE FUNCTION foo(f1 int) RETURNS int AS $$
DECLARE
    f1 int;
BEGIN
    RETURN f1;
END
$$ LANGUAGE plpgsql;
WARNING:  variable "f1" shadows a previously defined variable
LINE 3: f1 int;

CREATE FUNCTION
```

### 43.13. Porting from Oracle PL/SQL

This section explains differences between PostgreSQL's PL/pgSQL language and Oracle's PL/SQL language, to help developers who port applications from Oracle® to PostgreSQL.

PL/pgSQL is similar to PL/SQL in many aspects. It is a block-structured, imperative language, and all variables have to be declared. Assignments, loops, conditionals are similar. The main differences you should keep in mind when porting from PL/SQL to PL/pgSQL are:

- If a name used in a SQL command could be either a column name of a table or a reference to a variable of the function, PL/SQL treats it as a column name. This corresponds to PL/pgSQL's `plpgsql.variable_conflict = use_column` behavior, which is not the default, as ex-
plained in Section 43.11.1. It’s often best to avoid such ambiguities in the first place, but if you have to port a large amount of code that depends on this behavior, setting variable_conflict may be the best solution.

• In PostgreSQL the function body must be written as a string literal. Therefore you need to use dollar quoting or escape single quotes in the function body. (See Section 43.12.1.)

• Data type names often need translation. For example, in Oracle string values are commonly declared as being of type varchar2, which is a non-SQL-standard type. In PostgreSQL, use type varchar or text instead. Similarly, replace type number with numeric, or use some other numeric data type if there’s a more appropriate one.

• Instead of packages, use schemas to organize your functions into groups.

• Since there are no packages, there are no package-level variables either. This is somewhat annoying. You can keep per-session state in temporary tables instead.

• Integer FOR loops with REVERSE work differently: PL/SQL counts down from the second number to the first, while PL/pgSQL counts down from the first number to the second, requiring the loop bounds to be swapped when porting. This incompatibility is unfortunate but is unlikely to be changed. (See Section 43.6.5.5.)

• FOR loops over queries (other than cursors) also work differently: the target variable(s) must have been declared, whereas PL/SQL always declares them implicitly. An advantage of this is that the variable values are still accessible after the loop exits.

• There are various notational differences for the use of cursor variables.

43.13.1. Porting Examples

Example 43.9 shows how to port a simple function from PL/SQL to PL/pgSQL.

Example 43.9. Porting a Simple Function from PL/SQL to PL/pgSQL

Here is an Oracle PL/SQL function:

```sql
CREATE OR REPLACE FUNCTION cs_fmt_browser_version(v_name varchar2,
v_version varchar2)
RETURN varchar2 IS
BEGIN
  IF v_version IS NULL THEN
    RETURN v_name;
  END IF;
  RETURN v_name || '/' || v_version;
END;
/
show errors;
```

Let’s go through this function and see the differences compared to PL/pgSQL:

• The type name varchar2 has to be changed to varchar or text. In the examples in this section, we’ll use varchar, but text is often a better choice if you do not need specific string length limits.

• The RETURN key word in the function prototype (not the function body) becomes RETURNS in PostgreSQL. Also, IS becomes AS, and you need to add a LANGUAGE clause because PL/pgSQL is not the only possible function language.
PL/pgSQL - SQL Procedural Language

- In PostgreSQL, the function body is considered to be a string literal, so you need to use quote marks or dollar quotes around it. This substitutes for the terminating / in the Oracle approach.

- The show errors command does not exist in PostgreSQL, and is not needed since errors are reported automatically.

This is how this function would look when ported to PostgreSQL:

```sql
CREATE OR REPLACE FUNCTION cs_fmt_browser_version(v_name varchar, v_version varchar) RETURNS varchar AS $$
BEGIN
  IF v_version IS NULL THEN
    RETURN v_name;
  END IF;
  RETURN v_name || '/' || v_version;
END;
$$ LANGUAGE plpgsql;
```

Example 43.10 shows how to port a function that creates another function and how to handle the ensuing quoting problems.

**Example 43.10. Porting a Function that Creates Another Function from PL/SQL to PL/pgSQL**

The following procedure grabs rows from a SELECT statement and builds a large function with the results in IF statements, for the sake of efficiency.

This is the Oracle version:

```sql
CREATE OR REPLACE PROCEDURE cs_update_referrer_type_proc IS
  CURSOR referrer_keys IS
    SELECT * FROM cs_referrer_keys
    ORDER BY try_order;
  func_cmd VARCHAR(4000);
BEGIN
  func_cmd := 'CREATE OR REPLACE FUNCTION
    cs_find_referrer_type(v_host IN VARCHAR2,
      v_domain IN VARCHAR2, v_url IN VARCHAR2) RETURN
    VARCHAR2 IS BEGIN';
  FOR referrer_key IN referrer_keys LOOP
    func_cmd := func_cmd ||
      ' IF v_' || referrer_key.kind
      || ' LIKE ''' || referrer_key.key_string
      || ''' THEN RETURN ''' || referrer_key.referrer_type
      || ' ''; END IF;';
  END LOOP;
  func_cmd := func_cmd || ' RETURN NULL; END;';
  EXECUTE IMMEDIATE func_cmd;
END;
/
show errors;
```
Here is how this function would end up in PostgreSQL:

```sql
CREATE OR REPLACE PROCEDURE cs_update_referrer_type_proc() AS $func$
DECLARE
    referrer_keys CURSOR IS
        SELECT * FROM cs_referrer_keys
        ORDER BY try_order;
    func_body text;
    func_cmd text;
BEGIN
    func_body := 'BEGIN';
    FOR referrer_key IN referrer_keys LOOP
        func_body := func_body ||
            ' IF v_' || referrer_key.kind || ' LIKE ' || quote_literal(referrer_key.key_string)
            || ' THEN RETURN ' || quote_literal(referrer_key.referrer_type)
            || '; END IF; ';
    END LOOP;
    func_body := func_body || ' RETURN NULL; END;';

    func_cmd :=
        'CREATE OR REPLACE FUNCTION cs_find_referrer_type(v_host varchar,
                        v_domain varchar,
                        v_url varchar)
            RETURNS varchar AS '
            || quote_literal(func_body)
            || ' LANGUAGE plpgsql;' ;

    EXECUTE func_cmd;
END;
$func$
```

Notice how the body of the function is built separately and passed through `quote_literal` to double any quote marks in it. This technique is needed because we cannot safely use dollar quoting for defining the new function: we do not know for sure what strings will be interpolated from the `referrer_key.key_string` field. (We are assuming here that `referrer_key.kind` can be trusted to always be `host`, `domain`, or `url`, but `referrer_key.key_string` might be anything, in particular it might contain dollar signs.) This function is actually an improvement on the Oracle original, because it will not generate broken code when `referrer_key.key_string` or `referrer_key.referrer_type` contain quote marks.

Example 43.11 shows how to port a function with `OUT` parameters and string manipulation. PostgreSQL does not have a built-in `instr` function, but you can create one using a combination of other functions. In Section 43.13.3 there is a PL/pgSQL implementation of `instr` that you can use to make your porting easier.

**Example 43.11. Porting a Procedure With String Manipulation and OUT Parameters from PL/SQL to PL/pgSQL**

The following Oracle PL/SQL procedure is used to parse a URL and return several elements (host, path, and query).
CREATE OR REPLACE PROCEDURE cs_parse_url(
  v_url IN VARCHAR2,
  v_host OUT VARCHAR2,  -- This will be passed back
  v_path OUT VARCHAR2,  -- This one too
  v_query OUT VARCHAR2) -- And this one
IS
  a_pos1 INTEGER;
  a_pos2 INTEGER;
BEGIN
  v_host := NULL;
  v_path := NULL;
  v_query := NULL;
  a_pos1 := instr(v_url, '//');
  IF a_pos1 = 0 THEN
    RETURN;
  END IF;
  a_pos2 := instr(v_url, '/', a_pos1 + 2);
  IF a_pos2 = 0 THEN
    v_host := substr(v_url, a_pos1 + 2);
    v_path := '/';
    RETURN;
  END IF;
  v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
  a_pos1 := instr(v_url, '?', a_pos2 + 1);
  IF a_pos1 = 0 THEN
    v_path := substr(v_url, a_pos2);
    RETURN;
  END IF;
  v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
  v_query := substr(v_url, a_pos1 + 1);
END;
/
show errors;

Here is a possible translation into PL/pgSQL:

CREATE OR REPLACE FUNCTION cs_parse_url(
  v_url IN VARCHAR,
  v_host OUT VARCHAR,  -- This will be passed back
  v_path OUT VARCHAR,  -- This one too
  v_query OUT VARCHAR) -- And this one
AS $$
DECLARE
  a_pos1 INTEGER;
  a_pos2 INTEGER;
BEGIN
  v_host := NULL;
  v_path := NULL;
  v_query := NULL;
  a_pos1 := instr(v_url, '//');
  IF a_pos1 = 0 THEN
    RETURN;
  END IF;
  a_pos2 := instr(v_url, '/', a_pos1 + 2);
  IF a_pos2 = 0 THEN
    v_host := substr(v_url, a_pos1 + 2);
    v_path := '/';
    RETURN;
  END IF;
  v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
  a_pos1 := instr(v_url, '?', a_pos2 + 1);
  IF a_pos1 = 0 THEN
    v_path := substr(v_url, a_pos2);
    RETURN;
  END IF;
  v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
  v_query := substr(v_url, a_pos1 + 1);
END;
/
show errors;
IF a_pos1 = 0 THEN
    RETURN;
END IF;

a_pos2 := instr(v_url, '/', a_pos1 + 2);
IF a_pos2 = 0 THEN
    v_host := substr(v_url, a_pos1 + 2);
    v_path := '/';
    RETURN;
END IF;

v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
a_pos1 := instr(v_url, '?', a_pos2 + 1);

IF a_pos1 = 0 THEN
    v_path := substr(v_url, a_pos2);
    RETURN;
END IF;

v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
    v_query := substr(v_url, a_pos1 + 1);
END;
$$ LANGUAGE plpgsql;

This function could be used like this:

SELECT * FROM cs_parse_url('http://foobar.com/query.cgi?baz');

Example 43.12 shows how to port a procedure that uses numerous features that are specific to Oracle.

**Example 43.12. Porting a Procedure from PL/SQL to PL/pgSQL**

The Oracle version:

CREATE OR REPLACE PROCEDURE cs_create_job(v_job_id IN INTEGER) IS
    a_running_job_count INTEGER;
BEGIN
    LOCK TABLE cs_jobs IN EXCLUSIVE MODE;

    SELECT count(*) INTO a_running_job_count FROM cs_jobs WHERE
    end_stamp IS NULL;

    IF a_running_job_count > 0 THEN
        COMMIT; -- free lock
        raise_application_error(-20000,
            'Unable to create a new job: a job is currently
            running.');
    END IF;

    DELETE FROM cs_active_job;
    INSERT INTO cs_active_job(job_id) VALUES (v_job_id);

    BEGIN
        INSERT INTO cs_jobs (job_id, start_stamp) VALUES (v_job_id, sysdate);
        EXCEPTION
            WHEN dup_val_on_index THEN NULL; -- don't worry if it
            already exists
This is how we could port this procedure to PL/pgSQL:

```sql
CREATE OR REPLACE PROCEDURE cs_create_job(v_job_id integer) AS $$
DECLARE
  a_running_job_count integer;
BEGIN
  LOCK TABLE cs_jobs IN EXCLUSIVE MODE;
  SELECT count(*) INTO a_running_job_count FROM cs_jobs WHERE end_stamp IS NULL;
  IF a_running_job_count > 0 THEN
    COMMIT; -- free lock
    RAISE EXCEPTION 'Unable to create a new job: a job is currently running'; -- 1
  END IF;

  DELETE FROM cs_active_job;
  INSERT INTO cs_active_job(job_id) VALUES (v_job_id);
  BEGIN
    INSERT INTO cs_jobs (job_id, start_stamp) VALUES (v_job_id, now());
    EXCEPTION
    WHEN unique_violation THEN -- 2
      -- don't worry if it already exists
  END;
  COMMIT;
END;
$$ LANGUAGE plpgsql;
```

1 The syntax of RAISE is considerably different from Oracle's statement, although the basic case `RAISE exception_name` works similarly.
2 The exception names supported by PL/pgSQL are different from Oracle's. The set of built-in exception names is much larger (see Appendix A). There is not currently a way to declare user-defined exception names, although you can throw user-chosen SQLSTATE values instead.

### 43.13.2. Other Things to Watch For

This section explains a few other things to watch for when porting Oracle PL/SQL functions to PostgreSQL.

#### 43.13.2.1. Implicit Rollback after Exceptions

In PL/pgSQL, when an exception is caught by an `EXCEPTION` clause, all database changes since the block's `BEGIN` are automatically rolled back. That is, the behavior is equivalent to what you'd get in Oracle with:

```sql
BEGIN
  SAVEPOINT s1;
  ... code here ...
```

1220
EXCEPTION
  WHEN ... THEN
    ROLLBACK TO s1;
    ... code here ...
  WHEN ... THEN
    ROLLBACK TO s1;
    ... code here ...
END;

If you are translating an Oracle procedure that uses SAVEPOINT and ROLLBACK TO in this style, your task is easy: just omit the SAVEPOINT and ROLLBACK TO. If you have a procedure that uses SAVEPOINT and ROLLBACK TO in a different way then some actual thought will be required.

43.13.2.2. EXECUTE

The PL/pgSQL version of EXECUTE works similarly to the PL/SQL version, but you have to remember to use quote_literal and quote_ident as described in Section 43.5.4. Constructs of the type EXECUTE 'SELECT * FROM $1'; will not work reliably unless you use these functions.

43.13.2.3. Optimizing PL/pgSQL Functions

PostgreSQL gives you two function creation modifiers to optimize execution: “volatility” (whether the function always returns the same result when given the same arguments) and “strictness” (whether the function returns null if any argument is null). Consult the CREATE FUNCTION reference page for details.

When making use of these optimization attributes, your CREATE FUNCTION statement might look something like this:

CREATE FUNCTION foo(...) RETURNS integer AS $$
...$$ LANGUAGE plpgsql STRICT IMMUTABLE;

43.13.3. Appendix

This section contains the code for a set of Oracle-compatible instr functions that you can use to simplify your porting efforts.

--
-- instr functions that mimic Oracle's counterpart
-- Syntax: instr(string1, string2 [, n [, m]])
-- where [] denotes optional parameters.
--
-- Search string1, beginning at the nth character, for the mth occurrence
-- of string2. If n is negative, search backwards, starting at the abs(n)'th
-- character from the end of string1.
-- If n is not passed, assume 1 (search starts at first character).
-- If m is not passed, assume 1 (find first occurrence).
-- Returns starting index of string2 in string1, or 0 if string2 is not found.
--
CREATE FUNCTION instr(varchar, varchar) RETURNS integer AS $$
BEGIN
  RETURN instr($1, $2, 1);
$$
CREATE FUNCTION instr(string varchar, string_to_search_for varchar,
  beg_index integer)
RETURNS integer AS $$
DECLARE
  pos integer NOT NULL DEFAULT 0;
  temp_str varchar;
  beg integer;
  length integer;
  ss_length integer;
BEGIN
  IF beg_index > 0 THEN
    temp_str := substring(string FROM beg_index);
    pos := position(string_to_search_for IN temp_str);
    IF pos = 0 THEN
      RETURN 0;
    ELSE
      RETURN pos + beg_index - 1;
    END IF;
  ELSIF beg_index < 0 THEN
    ss_length := char_length(string_to_search_for);
    length := char_length(string);
    beg := length + 1 + beg_index;
    WHILE beg > 0 LOOP
      temp_str := substring(string FROM beg FOR ss_length);
      IF string_to_search_for = temp_str THEN
        RETURN beg;
      END IF;
      beg := beg - 1;
    END LOOP;
    RETURN 0;
  ELSE
    RETURN 0;
  END IF;
END;
$$ LANGUAGE plpgsql STRICT IMMUTABLE;

CREATE FUNCTION instr(string varchar, string_to_search_for varchar,
  beg_index integer, occur_index integer)
RETURNS integer AS $$
DECLARE
  pos integer NOT NULL DEFAULT 0;
  occur_number integer NOT NULL DEFAULT 0;
  temp_str varchar;
  beg integer;
  i integer;
  length integer;
  ss_length integer;
BEGIN
  IF occur_index <= 0 THEN

RAISE 'argument ''%'' is out of range', occur_index
   USING ERRCODE = '22003';
END IF;

IF beg_index > 0 THEN
   beg := beg_index - 1;
   FOR i IN 1..occur_index LOOP
      temp_str := substring(string FROM beg + 1);
      pos := position(string_to_search_for IN temp_str);
      IF pos = 0 THEN
         RETURN 0;
      END IF;
      beg := beg + pos;
   END LOOP;
   RETURN beg;
ELSIF beg_index < 0 THEN
   ss_length := char_length(string_to_search_for);
   length := char_length(string);
   beg := length + 1 + beg_index;
   WHILE beg > 0 LOOP
      temp_str := substring(string FROM beg FOR ss_length);
      IF string_to_search_for = temp_str THEN
         occur_number := occur_number + 1;
         IF occur_number = occur_index THEN
            RETURN beg;
         END IF;
      END IF;
      beg := beg - 1;
   END LOOP;
   RETURN 0;
ELSE
   RETURN 0;
END IF;
END;
$$ LANGUAGE plpgsql STRICT IMMUTABLE;
Chapter 44. PL/Tcl - Tcl Procedural Language

PL/Tcl is a loadable procedural language for the PostgreSQL database system that enables the Tcl language\(^1\) to be used to write PostgreSQL functions.

44.1. Overview

PL/Tcl offers most of the capabilities a function writer has in the C language, with a few restrictions, and with the addition of the powerful string processing libraries that are available for Tcl.

One compelling good restriction is that everything is executed from within the safety of the context of a Tcl interpreter. In addition to the limited command set of safe Tcl, only a few commands are available to access the database via SPI and to raise messages via `elog()`. PL/Tcl provides no way to access internals of the database server or to gain OS-level access under the permissions of the PostgreSQL server process, as a C function can do. Thus, unprivileged database users can be trusted to use this language; it does not give them unlimited authority.

The other notable implementation restriction is that Tcl functions cannot be used to create input/output functions for new data types.

Sometimes it is desirable to write Tcl functions that are not restricted to safe Tcl. For example, one might want a Tcl function that sends email. To handle these cases, there is a variant of PL/Tcl called PL/TclU (for untrusted Tcl). This is exactly the same language except that a full Tcl interpreter is used. If PL/TclU is used, it must be installed as an untrusted procedural language so that only database superusers can create functions in it. The writer of a PL/TclU function must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator.

The shared object code for the PL/Tcl and PL/TclU call handlers is automatically built and installed in the PostgreSQL library directory if Tcl support is specified in the configuration step of the installation procedure. To install PL/Tcl and/or PL/TclU in a particular database, use the `CREATE EXTENSION` command, for example `CREATE EXTENSION pltcl` or `CREATE EXTENSION pltclu`.

44.2. PL/Tcl Functions and Arguments

To create a function in the PL/Tcl language, use the standard CREATE FUNCTION syntax:

```sql
CREATE FUNCTION funcname (argument-types) RETURNS return-type AS $$
    # PL/Tcl function body
    $$ LANGUAGE pltcl;
```

PL/TclU is the same, except that the language has to be specified as `pltclu`.

The body of the function is simply a piece of Tcl script. When the function is called, the argument values are passed to the Tcl script as variables named 1 ... n. The result is returned from the Tcl code in the usual way, with a `return` statement. In a procedure, the return value from the Tcl code is ignored.

For example, a function returning the greater of two integer values could be defined as:

```sql
CREATE FUNCTION tcl_max(integer, integer) RETURNS integer AS $$
    if ($1 > $2) {return $1}
    $$ LANGUAGE pltcl;
```

---

\(^1\) [http://www.tcl.tk/](http://www.tcl.tk/)
Note the clause STRICT, which saves us from having to think about null input values: if a null value is passed, the function will not be called at all, but will just return a null result automatically.

In a nonstrict function, if the actual value of an argument is null, the corresponding \$n variable will be set to an empty string. To detect whether a particular argument is null, use the function argisnull. For example, suppose that we wanted tcl_max with one null and one nonnull argument to return the nonnull argument, rather than null:

```
CREATE FUNCTION tcl_max(integer, integer) RETURNS integer AS $$
    if {[argisnull 1]} {
        if {[argisnull 2]} { return_null }
        return $2
    }
    if {[argisnull 2]} { return $1 }
    if ($1 > $2) { return $1 }
    return $2
$$ LANGUAGE pltcl;
```

As shown above, to return a null value from a PL/Tcl function, execute return_null. This can be done whether the function is strict or not.

Composite-type arguments are passed to the function as Tcl arrays. The element names of the array are the attribute names of the composite type. If an attribute in the passed row has the null value, it will not appear in the array. Here is an example:

```
CREATE TABLE employee (
    name text,
    salary integer,
    age integer
);

CREATE FUNCTION overpaid(employee) RETURNS boolean AS $$
    if {200000.0 < $1(salary)} {
        return "t"
    }
    if {$1(age) < 30 && 100000.0 < $1(salary)} {
        return "t"
    }
    return "f"
$$ LANGUAGE pltcl;
```

As shown above, to return a null value from a PL/Tcl function, execute return_null. This can be done whether the function is strict or not.

Composite-type arguments are passed to the function as Tcl arrays. The element names of the array are the attribute names of the composite type. If an attribute in the passed row has the null value, it will not appear in the array. Here is an example:

```
CREATE TABLE employee (
    name text,
    salary integer,
    age integer
);

CREATE FUNCTION overpaid(employee) RETURNS boolean AS $$
    if {200000.0 < $1(salary)} {
        return "t"
    }
    if {$1(age) < 30 && 100000.0 < $1(salary)} {
        return "t"
    }
    return "f"
$$ LANGUAGE pltcl;
```

PL/Tcl functions can return composite-type results, too. To do this, the Tcl code must return a list of column name/value pairs matching the expected result type. Any column names omitted from the list are returned as nulls, and an error is raised if there are unexpected column names. Here is an example:

```
CREATE FUNCTION square_cube(in int, out squared int, out cubed int) AS $$
    return [list squared [expr {$1 * $1}] cubed [expr {$1 * $1 * $1}]]
$$ LANGUAGE pltcl;
```

Output arguments of procedures are returned in the same way, for example:
CREATE PROCEDURE tcl_triple(INOUT a integer, INOUT b integer) AS $$
return [list a [expr {$1 * 3}] b [expr {$2 * 3}]]
$$ LANGUAGE pltcl;

CALL tcl_triple(5, 10);

Tip

The result list can be made from an array representation of the desired tuple with the array get Tcl command. For example:

CREATE FUNCTION raise_pay(employee, delta int) RETURNS employee AS $$
set l(salary) [expr {$1(salary) + $2}]
return [array get l]
$$ LANGUAGE pltcl;

PL/Tcl functions can return sets. To do this, the Tcl code should call return_next once per row to be returned, passing either the appropriate value when returning a scalar type, or a list of column name/value pairs when returning a composite type. Here is an example returning a scalar type:

CREATE FUNCTION sequence(int, int) RETURNS SETOF int AS $$
for {set i $1} {$i < $2} {incr i} {
    return_next $i
}
$$ LANGUAGE pltcl;

and here is one returning a composite type:

CREATE FUNCTION table_of_squares(int, int) RETURNS TABLE (x int, x2 int) AS $$
for {set i $1} {$i < $2} {incr i} {
    return_next [list x $i x2 [expr {$i * $i}]]
}
$$ LANGUAGE pltcl;

44.3. Data Values in PL/Tcl

The argument values supplied to a PL/Tcl function's code are simply the input arguments converted to text form (just as if they had been displayed by a SELECT statement). Conversely, the return and return_next commands will accept any string that is acceptable input format for the function's declared result type, or for the specified column of a composite result type.

44.4. Global Data in PL/Tcl

Sometimes it is useful to have some global data that is held between two calls to a function or is shared between different functions. This is easily done in PL/Tcl, but there are some restrictions that must be understood.

For security reasons, PL/Tcl executes functions called by any one SQL role in a separate Tcl interpreter for that role. This prevents accidental or malicious interference by one user with the behavior of another user's PL/Tcl functions. Each such interpreter will have its own values for any “global” Tcl variables. Thus, two PL/Tcl functions will share the same global variables if and only if they are executed by the same user.
same SQL role. In an application wherein a single session executes code under multiple SQL roles (via SECURITY DEFINER functions, use of SET ROLE, etc.) you may need to take explicit steps to ensure that PL/Tcl functions can share data. To do that, make sure that functions that should communicate are owned by the same user, and mark them SECURITY DEFINER. You must of course take care that such functions can't be used to do anything unintended.

All PL/TcLU functions used in a session execute in the same Tcl interpreter, which of course is distinct from the interpreter(s) used for PL/Tcl functions. So global data is automatically shared between PL/TcLU functions. This is not considered a security risk because all PL/TcLU functions execute at the same trust level, namely that of a database superuser.

To help protect PL/Tcl functions from unintentionally interfering with each other, a global array is made available to each function via the upvar command. The global name of this variable is the function's internal name, and the local name is GD. It is recommended that GD be used for persistent private data of a function. Use regular Tcl global variables only for values that you specifically intend to be shared among multiple functions. (Note that the GD arrays are only global within a particular interpreter, so they do not bypass the security restrictions mentioned above.)

An example of using GD appears in the spi_execp example below.

### 44.5. Database Access from PL/Tcl

The following commands are available to access the database from the body of a PL/Tcl function:

```
spi_exec ?-count n? ?-array name? command ?loop-body?
```

- **spi_exec** executes an SQL command given as a string. An error in the command causes an error to be raised. Otherwise, the return value of `spi_exec` is the number of rows processed (selected, inserted, updated, or deleted) by the command, or zero if the command is a utility statement. In addition, if the command is a SELECT statement, the values of the selected columns are placed in Tcl variables as described below.

- The optional -count value tells `spi_exec` the maximum number of rows to process in the command. The effect of this is comparable to setting up a query as a cursor and then saying **FETCH n**.

- If the command is a SELECT statement, the values of the result columns are placed into Tcl variables named after the columns. If the -array option is given, the column values are instead stored into elements of the named associative array, with the column names used as array indexes. In addition, the current row number within the result (counting from zero) is stored into the array element named ".tupno", unless that name is in use as a column name in the result.

- If the command is a SELECT statement and no loop-body script is given, then only the first row of results are stored into Tcl variables or array elements; remaining rows, if any, are ignored. No storing occurs if the query returns no rows. (This case can be detected by checking the result of `spi_exec`.) For example:

```
spi_exec "SELECT count(*) AS cnt FROM pg_proc"
```

will set the Tcl variable $cnt to the number of rows in the pg_proc system catalog.

- If the optional loop-body argument is given, it is a piece of Tcl script that is executed once for each row in the query result. (loop-body is ignored if the given command is not a SELECT.) The values of the current row's columns are stored into Tcl variables or array elements before each iteration. For example:

```
spi_exec -array C "SELECT * FROM pg_class" {
```
elog DEBUG "have table $C(rename)"
}

will print a log message for every row of pg_class. This feature works similarly to other Tcl looping constructs; in particular continue and break work in the usual way inside the loop body.

If a column of a query result is null, the target variable for it is “unset” rather than being set.

```tcl
spi_prepare query typelist
```

Prepares and saves a query plan for later execution. The saved plan will be retained for the life of the current session.

The query can use parameters, that is, placeholders for values to be supplied whenever the plan is actually executed. In the query string, refer to parameters by the symbols $1 ... $n. If the query uses parameters, the names of the parameter types must be given as a Tcl list. (Write an empty list for typelist if no parameters are used.)

The return value from spi_prepare is a query ID to be used in subsequent calls to spi_execp. See spi_execp for an example.

```tcl
spi_execp ?-count n ?-array name ?-nulls string? queryid ?value-list? ?loop-body?
```

Executes a query previously prepared with spi_prepare. queryid is the ID returned by spi_prepare. If the query references parameters, a value-list must be supplied. This is a Tcl list of actual values for the parameters. The list must be the same length as the parameter type list previously given to spi_prepare. Omit value-list if the query has no parameters.

The optional value for -nulls is a string of spaces and ‘n’ characters telling spi_execp which of the parameters are null values. If given, it must have exactly the same length as the value-list. If it is not given, all the parameter values are nonnull.

Except for the way in which the query and its parameters are specified, spi_execp works just like spi_exec. The -count, -array, and loop-body options are the same, and so is the result value.

Here’s an example of a PL/Tcl function using a prepared plan:

```tcl
CREATE FUNCTION t1_count(integer, integer) RETURNS integer AS $$
if {{[ info exists GD(plan) ]}} {
    # prepare the saved plan on the first call
    set GD(plan) [ spi_prepare \n        "SELECT count(*) AS cnt FROM t1 WHERE num >= \$1
            AND num <= \$2" \n        [ list int4 int4 ] ]
}
spi_execp -count 1 $GD(plan) [ list $1 $2 ]
return $cnt
$$ LANGUAGE pltcl;
```

We need backslashes inside the query string given to spi_prepare to ensure that the $n markers will be passed through to spi_prepare as-is, and not replaced by Tcl variable substitution.

```tcl
spi_lastoid
```

Returns the OID of the row inserted by the last spi_exec or spi_execp, if the command was a single-row INSERT and the modified table contained OIDs. (If not, you get zero.)
subtransaction command

The Tcl script contained in command is executed within a SQL subtransaction. If the script returns an error, that entire subtransaction is rolled back before returning the error out to the surrounding Tcl code. See Section 44.9 for more details and an example.

quote string

Doubles all occurrences of single quote and backslash characters in the given string. This can be used to safely quote strings that are to be inserted into SQL commands given to spi_exec or spi_prepare. For example, think about an SQL command string like:

"SELECT '$val' AS ret"

where the Tcl variable val actually contains doesn't. This would result in the final command string:

SELECT 'doesn't' AS ret

which would cause a parse error during spi_exec or spi_prepare. To work properly, the submitted command should contain:

SELECT 'doesn''t' AS ret

which can be formed in PL/Tcl using:

"SELECT '[ quote $val ]' AS ret"

One advantage of spi_execp is that you don't have to quote parameter values like this, since the parameters are never parsed as part of an SQL command string.

eol log level msg

Emits a log or error message. Possible levels are DEBUG, LOG, INFO, NOTICE, WARNING, ERROR, and FATAL. ERROR raises an error condition; if this is not trapped by the surrounding Tcl code, the error propagates out to the calling query, causing the current transaction or subtransaction to be aborted. This is effectively the same as the Tcl error command. FATAL aborts the transaction and causes the current session to shut down. (There is probably no good reason to use this error level in PL/Tcl functions, but it's provided for completeness.) The other levels only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the log_min_messages and client_min_messages configuration variables. See Chapter 19 and Section 44.8 for more information.

44.6. Trigger Functions in PL/Tcl

Trigger functions can be written in PL/Tcl. PostgreSQL requires that a function that is to be called as a trigger must be declared as a function with no arguments and a return type of trigger.

The information from the trigger manager is passed to the function body in the following variables:

$TG_name

The name of the trigger from the CREATE TRIGGER statement.

$TG_relid

The object ID of the table that caused the trigger function to be invoked.
$TG\_table\_name

The name of the table that caused the trigger function to be invoked.

$TG\_table\_schema

The schema of the table that caused the trigger function to be invoked.

$TG\_relatts

A Tcl list of the table column names, prefixed with an empty list element. So looking up a column name in the list with Tcl's \texttt{lsearch} command returns the element's number starting with 1 for the first column, the same way the columns are customarily numbered in PostgreSQL. (Empty list elements also appear in the positions of columns that have been dropped, so that the attribute numbering is correct for columns to their right.)

$TG\_when

The string \texttt{BEFORE}, \texttt{AFTER}, or \texttt{INSTEAD OF}, depending on the type of trigger event.

$TG\_level

The string \texttt{ROW} or \texttt{STATEMENT} depending on the type of trigger event.

$TG\_op

The string \texttt{INSERT}, \texttt{UPDATE}, \texttt{DELETE}, or \texttt{TRUNCATE} depending on the type of trigger event.

$NEW

An associative array containing the values of the new table row for \texttt{INSERT} or \texttt{UPDATE} actions, or empty for \texttt{DELETE}. The array is indexed by column name. Columns that are null will not appear in the array. This is not set for statement-level triggers.

$OLD

An associative array containing the values of the old table row for \texttt{UPDATE} or \texttt{DELETE} actions, or empty for \texttt{INSERT}. The array is indexed by column name. Columns that are null will not appear in the array. This is not set for statement-level triggers.

$\texttt{args}$

A Tcl list of the arguments to the function as given in the \texttt{CREATE TRIGGER} statement. These arguments are also accessible as $1 \ldots n$ in the function body.

The return value from a trigger function can be one of the strings \texttt{OK} or \texttt{SKIP}, or a list of column name/value pairs. If the return value is \texttt{OK}, the operation (\texttt{INSERT/UPDATE/DELETE}) that fired the trigger will proceed normally. \texttt{SKIP} tells the trigger manager to silently suppress the operation for this row. If a list is returned, it tells PL/Tcl to return a modified row to the trigger manager; the contents of the modified row are specified by the column names and values in the list. Any columns not mentioned in the list are set to null. Returning a modified row is only meaningful for row-level \texttt{BEFORE INSERT} or \texttt{UPDATE} triggers, for which the modified row will be inserted instead of the one given in $\texttt{NEW}$; or for row-level \texttt{INSTEAD OF INSERT} or \texttt{UPDATE} triggers where the returned row is used as the source data for \texttt{INSERT RETURNING} or \texttt{UPDATE RETURNING} clauses. In row-level \texttt{BEFORE DELETE} or \texttt{INSTEAD OF DELETE} triggers, returning a modified row has the same effect as returning \texttt{OK}, that is the operation proceeds. The trigger return value is ignored for all other types of triggers.

<table>
<thead>
<tr>
<th>Tip</th>
</tr>
</thead>
<tbody>
<tr>
<td>The result list can be made from an array representation of the modified tuple with the \texttt{array get} Tcl command.</td>
</tr>
</tbody>
</table>

1230
Here's a little example trigger function that forces an integer value in a table to keep track of the number of updates that are performed on the row. For new rows inserted, the value is initialized to 0 and then incremented on every update operation.

```pltcl
CREATE FUNCTION trigfunc_modcount() RETURNS trigger AS $$
switch $TG_op {
    INSERT {
       set NEW($1) 0
    }
    UPDATE {
       set NEW($1) $OLD($1)
       incr NEW($1)
    }
    default {
       return OK
    }
}
return [array get NEW]
$$ LANGUAGE pltcl;
CREATE TABLE mytab (num integer, description text, modcnt integer);
CREATE TRIGGER trig_mytab_modcount BEFORE INSERT OR UPDATE ON mytab
FOR EACH ROW EXECUTE FUNCTION trigfunc_modcount('modcnt');
Notice that the trigger function itself does not know the column name; that's supplied from the trigger arguments. This lets the trigger function be reused with different tables.

### 44.7. Event Trigger Functions in PL/Tcl

Event trigger functions can be written in PL/Tcl. PostgreSQL requires that a function that is to be called as an event trigger must be declared as a function with no arguments and a return type of event_trigger.

The information from the trigger manager is passed to the function body in the following variables:

- `$TG_event`
  - The name of the event the trigger is fired for.
- `$TG_tag`
  - The command tag for which the trigger is fired.

The return value of the trigger function is ignored.

Here's a little example event trigger function that simply raises a NOTICE message each time a supported command is executed:

```pltcl
CREATE OR REPLACE FUNCTION tclsnitch() RETURNS event_trigger AS $$
   elog NOTICE "tclsnitch: $TG_event $TG_tag"
$$ LANGUAGE pltcl;
CREATE EVENT TRIGGER tcl_a_snitch ON ddl_command_start EXECUTE FUNCTION tclsnitch();
```

### 44.8. Error Handling in PL/Tcl

1231
PL/Tcl - Tcl Procedural Language

Tcl code within or called from a PL/Tcl function can raise an error, either by executing some invalid operation or by generating an error using the Tcl error command or PL/Tcl's elog command. Such errors can be caught within Tcl using the Tcl catch command. If an error is not caught but is allowed to propagate out to the top level of execution of the PL/Tcl function, it is reported as a SQL error in the function's calling query.

Conversely, SQL errors that occur within PL/Tcl's spi_exec, spi_prepare, and spi_execp commands are reported as Tcl errors, so they are catchable by Tcl's catch command. (Each of these PL/Tcl commands runs its SQL operation in a subtransaction, which is rolled back on error, so that any partially-completed operation is automatically cleaned up.) Again, if an error propagates out to the top level without being caught, it turns back into a SQL error.

Tcl provides an errorCode variable that can represent additional information about an error in a form that is easy for Tcl programs to interpret. The contents are in Tcl list format, and the first word identifies the subsystem or library reporting the error; beyond that the contents are left to the individual subsystem or library. For database errors reported by PL/Tcl commands, the first word is POSTGRES, the second word is the PostgreSQL version number, and additional words are field name/value pairs providing detailed information about the error. Fields SQLSTATE, condition, and message are always supplied (the first two represent the error code and condition name as shown in Appendix A). Fields that may be present include detail, hint, context, schema, table, column, datatype, constraint, statement, cursor_position, filename, lineno, and funcname.

A convenient way to work with PL/Tcl's errorCode information is to load it into an array, so that the field names become array subscripts. Code for doing that might look like

```tcl
if {{catch { spi_exec $sql_command }}} {
    if {[lindex $::errorCode 0] == "POSTGRES"} {
        array set errorArray $::errorCode
        if {$errorArray(condition) == "undefined_table"} {
            # deal with missing table
        } else {
            # deal with some other type of SQL error
        }
    }
}
```

(The double colons explicitly specify that errorCode is a global variable.)

### 44.9. Explicit Subtransactions in PL/Tcl

Recovering from errors caused by database access as described in Section 44.8 can lead to an undesirable situation where some operations succeed before one of them fails, and after recovering from that error the data is left in an inconsistent state. PL/Tcl offers a solution to this problem in the form of explicit subtransactions.

Consider a function that implements a transfer between two accounts:

```tcl
CREATE FUNCTION transfer_funds() RETURNS void AS $$
if [catch {
    spi_exec "UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'"
    spi_exec "UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'"
} errmsg] {
    set result [format "error transferring funds: %s" $errmsg]
$$
```
If the second UPDATE statement results in an exception being raised, this function will log the failure, but the result of the first UPDATE will nevertheless be committed. In other words, the funds will be withdrawn from Joe's account, but will not be transferred to Mary's account. This happens because each spi_exec is a separate subtransaction, and only one of those subtransactions got rolled back.

To handle such cases, you can wrap multiple database operations in an explicit subtransaction, which will succeed or roll back as a whole. PL/Tcl provides a subtransaction command to manage this. We can rewrite our function as:

```tcl
CREATE FUNCTION transfer_funds2() RETURNS void AS $$
    if [catch {
        subtransaction {
            spi_exec "UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'"
            spi_exec "UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'"
        }
    } errormsg] {
        set result [format "error transferring funds: %s" $errormsg]
    } else {
        set result "funds transferred successfully"
    }
    spi_exec "INSERT INTO operations (result) VALUES ('\[quote $result\])'"
$$ LANGUAGE pltcl;
```

Note that use of catch is still required for this purpose. Otherwise the error would propagate to the top level of the function, preventing the desired insertion into the operations table. The subtransaction command does not trap errors, it only assures that all database operations executed inside its scope will be rolled back together when an error is reported.

A rollback of an explicit subtransaction occurs on any error reported by the contained Tcl code, not only errors originating from database access. Thus a regular Tcl exception raised inside a subtransaction command will also cause the subtransaction to be rolled back. However, non-error exits out of the contained Tcl code (for instance, due to return) do not cause a rollback.

### 44.10. Transaction Management

In a procedure called from the top level or an anonymous code block (DO command) called from the top level it is possible to control transactions. To commit the current transaction, call the commit command. To roll back the current transaction, call the rollback command. (Note that it is not possible to run the SQL commands COMMIT or ROLLBACK via spi_exec or similar. It has to be done using these functions.) After a transaction is ended, a new transaction is automatically started, so there is no separate command for that.

Here is an example:

```tcl
CREATE PROCEDURE transaction_test1()
LANGUAGE pltcl
```
AS $$
for {set i 0} {$i < 10} {incr i} {
    spi_exec "INSERT INTO test1 (a) VALUES ($i)"
    if {($i % 2 == 0)} {
        commit
    } else {
        rollback
    }
}
$$;

CALL transaction_test1();

Transactions cannot be ended when an explicit subtransaction is active.

44.11. PL/Tcl Configuration

This section lists configuration parameters that affect PL/Tcl.

pltcl.start_proc(string)

This parameter, if set to a nonempty string, specifies the name (possibly schema-qualified) of a parameterless PL/Tcl function that is to be executed whenever a new Tcl interpreter is created for PL/Tcl. Such a function can perform per-session initialization, such as loading additional Tcl code. A new Tcl interpreter is created when a PL/Tcl function is first executed in a database session, or when an additional interpreter has to be created because a PL/Tcl function is called by a new SQL role.

The referenced function must be written in the pltcl language, and must not be marked SECURITY DEFINER. (These restrictions ensure that it runs in the interpreter it's supposed to initialize.) The current user must have permission to call it, too.

If the function fails with an error it will abort the function call that caused the new interpreter to be created and propagate out to the calling query, causing the current transaction or subtransaction to be aborted. Any actions already done within Tcl won't be undone; however, that interpreter won't be used again. If the language is used again the initialization will be attempted again within a fresh Tcl interpreter.

Only superusers can change this setting. Although this setting can be changed within a session, such changes will not affect Tcl interpreters that have already been created.

pltclu.start_proc(string)

This parameter is exactly like pltcl.start_proc, except that it applies to PL/TclU. The referenced function must be written in the pltclu language.

44.12. Tcl Procedure Names

In PostgreSQL, the same function name can be used for different function definitions as long as the number of arguments or their types differ. Tcl, however, requires all procedure names to be distinct. PL/Tcl deals with this by making the internal Tcl procedure names contain the object ID of the function from the system table pg_proc as part of their name. Thus, PostgreSQL functions with the same name and different argument types will be different Tcl procedures, too. This is not normally a concern for a PL/Tcl programmer, but it might be visible when debugging.
Chapter 45. PL/Perl - Perl Procedural Language

PL/Perl is a loadable procedural language that enables you to write PostgreSQL functions in the Perl programming language.

The main advantage to using PL/Perl is that this allows use, within stored functions, of the manyfold “string munging” operators and functions available for Perl. Parsing complex strings might be easier using Perl than it is with the string functions and control structures provided in PL/pgSQL.

To install PL/Perl in a particular database, use CREATE EXTENSION plperl.

Tip
If a language is installed into template1, all subsequently created databases will have the language installed automatically.

Note
Users of source packages must specially enable the build of PL/Perl during the installation process. (Refer to Chapter 16 for more information.) Users of binary packages might find PL/Perl in a separate subpackage.

45.1. PL/Perl Functions and Arguments

To create a function in the PL/Perl language, use the standard CREATE FUNCTION syntax:

CREATE FUNCTION funcname (argument-types) RETURNS return-type AS $$
    # PL/Perl function body
$$ LANGUAGE plperl;

The body of the function is ordinary Perl code. In fact, the PL/Perl glue code wraps it inside a Perl subroutine. A PL/Perl function is called in a scalar context, so it can't return a list. You can return non-scalar values (arrays, records, and sets) by returning a reference, as discussed below.

In a PL/Perl procedure, any return value from the Perl code is ignored.

PL/Perl also supports anonymous code blocks called with the DO statement:

DO $$
    # PL/Perl code
$$ LANGUAGE plperl;

An anonymous code block receives no arguments, and whatever value it might return is discarded. Otherwise it behaves just like a function.

Note
The use of named nested subroutines is dangerous in Perl, especially if they refer to lexical variables in the enclosing scope. Because a PL/Perl function is wrapped in a

1 http://www.perl.org
subroutine, any named subroutine you place inside one will be nested. In general, it
is far safer to create anonymous subroutines which you call via a coderef. For more
information, see the entries for \texttt{Variable "\texttt{\%s}" will not stay shared and
Variable "\texttt{\%s}" is not available} in the \texttt{perldiag} man page, or search the
Internet for “perl nested named subroutine”.

The syntax of the \texttt{CREATE FUNCTION} command requires the function body to be written as a string
constant. It is usually most convenient to use dollar quoting (see Section 4.1.2.4) for the string constant.
If you choose to use escape string syntax \texttt{E"\texttt{\textbackslash\textbackslash}"}, you must double any single quote marks (')
and backslashes (\) used in the body of the function (see Section 4.1.2.1).

Arguments and results are handled as in any other Perl subroutine: arguments are passed in \texttt{@_}
and a
result value is returned with \texttt{return} or as the last expression evaluated in the function.

For example, a function returning the greater of two integer values could be defined as:

\begin{verbatim}
CREATE FUNCTION perl_max (integer, integer) RETURNS integer AS $$
  if ($_[0] > $_[1]) { return $_[0]; } 
  return $_[1];
$$ LANGUAGE plperl;
\end{verbatim}

\textbf{Note}

Arguments will be converted from the database's encoding to UTF-8 for use inside PL/Perl, and then converted from UTF-8 back to the database encoding upon return.

If an SQL null value is passed to a function, the argument value will appear as “undefined” in Perl.
The above function definition will not behave very nicely with null inputs (in fact, it will act as though
they are zeroes). We could add \texttt{STRICT} to the function definition to make PostgreSQL do something
more reasonable: if a null value is passed, the function will not be called at all, but will just return a
null result automatically. Alternatively, we could check for undefined inputs in the function body. For
example, suppose that we wanted \texttt{perl_max} with one null and one nonnull argument to return the
nonnull argument, rather than a null value:

\begin{verbatim}
CREATE FUNCTION perl_max (integer, integer) RETURNS integer AS $$
  my ($x, $y) = @_; 
  if (not defined $x) { 
    return undef if not defined $y; 
    return $y; 
  } 
  return $x if not defined $y; 
  return $x if $x > $y; 
  return $y;
$$ LANGUAGE plperl;
\end{verbatim}

As shown above, to return an SQL null value from a PL/Perl function, return an undefined value. This
can be done whether the function is strict or not.

Anything in a function argument that is not a reference is a string, which is in the standard PostgreSQL
external text representation for the relevant data type. In the case of ordinary numeric or text types, Perl
will just do the right thing and the programmer will normally not have to worry about it. However, in
other cases the argument will need to be converted into a form that is more usable in Perl. For example,
the \texttt{decode_bytea} function can be used to convert an argument of type \texttt{bytea} into unescaped
binary.
Similarly, values passed back to PostgreSQL must be in the external text representation format. For example, the `encode_bytea` function can be used to escape binary data for a return value of type `bytea`.

Perl can return PostgreSQL arrays as references to Perl arrays. Here is an example:

```plaintext
CREATE OR REPLACE function returns_array()
RETURNS text[] AS $$
  return ['a"b','c,d'],['e\f','g'];
$$ LANGUAGE plperl;
select returns_array();
```

Perl passes PostgreSQL arrays as a blessed `PostgreSQL::InServer::ARRAY` object. This object may be treated as an array reference or a string, allowing for backward compatibility with Perl code written for PostgreSQL versions below 9.1 to run. For example:

```plaintext
CREATE OR REPLACE FUNCTION concat_array_elements(text[]) RETURNS TEXT AS $$
  my $arg = shift;
  my $result = "";
  return undef if (!defined $arg);

  # as an array reference
  for (@$arg) {
    $result .= $_;
  }

  # also works as a string
  $result .= $arg;

  return $result;
$$ LANGUAGE plperl;
SELECT concat_array_elements(ARRAY['PL','/','Perl']);
```

Note

Multidimensional arrays are represented as references to lower-dimensional arrays of references in a way common to every Perl programmer.

Composite-type arguments are passed to the function as references to hashes. The keys of the hash are the attribute names of the composite type. Here is an example:

```plaintext
CREATE TABLE employee (  
  name text,  
  basesalary integer,  
  bonus integer
);
CREATE FUNCTION empcomp(employee) RETURNS integer AS $$
  my ($emp) = @_;  
  return $emp->{basesalary} + $emp->{bonus};
$$ LANGUAGE plperl;
```

1237
SELECT name, empcomp(employee.*) FROM employee;

A PL/Perl function can return a composite-type result using the same approach: return a reference to a hash that has the required attributes. For example:

CREATE TYPE testrowperl AS (f1 integer, f2 text, f3 text);

CREATE OR REPLACE FUNCTION perl_row() RETURNS testrowperl AS $$
return {f2 => 'hello', f1 => 1, f3 => 'world'};
$$ LANGUAGE plperl;

SELECT * FROM perl_row();

Any columns in the declared result data type that are not present in the hash will be returned as null values.

Similarly, output arguments of procedures can be returned as a hash reference:

CREATE PROCEDURE perl_triple(INOUT a integer, INOUT b integer) AS $$
my ($a, $b) = @_; return {a => $a * 3, b => $b * 3};
$$ LANGUAGE plperl;

CALL perl_triple(5, 10);

PL/Perl functions can also return sets of either scalar or composite types. Usually you'll want to return rows one at a time, both to speed up startup time and to keep from queuing up the entire result set in memory. You can do this with `return_next` as illustrated below. Note that after the last `return_next`, you must put either `return` or (better) `return undef`.

CREATE OR REPLACE FUNCTION perl_set_int(int) RETURNS SETOF INTEGER AS $$
foreach (0..$_[0]) {
    return_next($_);
}
return undef;
$$ LANGUAGE plperl;

SELECT * FROM perl_set_int(5);

CREATE OR REPLACE FUNCTION perl_set() RETURNS SETOF testrowperl AS $$
return_next({ f1 => 1, f2 => 'Hello', f3 => 'World' });
return_next({ f1 => 2, f2 => 'Hello', f3 => 'PostgreSQL' });
return_next({ f1 => 3, f2 => 'Hello', f3 => 'PL/Perl' });
return undef;
$$ LANGUAGE plperl;

For small result sets, you can return a reference to an array that contains either scalars, references to arrays, or references to hashes for simple types, array types, and composite types, respectively. Here are some simple examples of returning the entire result set as an array reference:

CREATE OR REPLACE FUNCTION perl_set_int(int) RETURNS SETOF INTEGER AS $$
return [];
$$
PL/Perl - Perl Procedural Language

```perlu
return [0..$_[0]]; $$ LANGUAGE plperl;
```

```sql
SELECT * FROM perl_set_int(5);
```

```perlu
CREATE OR REPLACE FUNCTION perl_set() RETURNS SETOF testrowperl AS $$
return [
    { f1 => 1, f2 => 'Hello', f3 => 'World' },
    { f1 => 2, f2 => 'Hello', f3 => 'PostgreSQL' },
    { f1 => 3, f2 => 'Hello', f3 => 'PL/Perl' }
]; $$ LANGUAGE plperl;
```

```sql
SELECT * FROM perl_set();
```

If you wish to use the strict pragma with your code you have a few options. For temporary global use you can SET plperl.use_strict to true. This will affect subsequent compilations of PL/Perl functions, but not functions already compiled in the current session. For permanent global use you can set plperl.use_strict to true in the postgresql.conf file.

For permanent use in specific functions you can simply put:

```perl
use strict;
```
at the top of the function body.

The feature pragma is also available to use if your Perl is version 5.10.0 or higher.

### 45.2. Data Values in PL/Perl

The argument values supplied to a PL/Perl function's code are simply the input arguments converted to text form (just as if they had been displayed by a SELECT statement). Conversely, the return and return_next commands will accept any string that is acceptable input format for the function's declared return type.

### 45.3. Built-in Functions

#### 45.3.1. Database Access from PL/Perl

Access to the database itself from your Perl function can be done via the following functions:

```
spi_exec_query (query [, max-rows])
```

`spi_exec_query` executes an SQL command and returns the entire row set as a reference to an array of hash references. You should only use this command when you know that the result set will be relatively small. Here is an example of a query (SELECT command) with the optional maximum number of rows:

```perl
$rv = spi_exec_query('SELECT * FROM my_table', 5);
```

This returns up to 5 rows from the table `my_table`. If `my_table` has a column `my_column`, you can get that value from row $i of the result like this:

```perl
$foo = $rv->{rows}[$i]->{my_column};
```
The total number of rows returned from a SELECT query can be accessed like this:

```
$nrows = $rv->{'processed'};
```

Here is an example using a different command type:

```
$query = "INSERT INTO my_table VALUES (1, 'test')";
$rv = spi_exec_query($query);
```

You can then access the command status (e.g., SPI_OK_INSERT) like this:

```
$res = $rv->{'status'};
```

To get the number of rows affected, do:

```
$nrows = $rv->{'processed'};
```

Here is a complete example:

```
CREATE TABLE test (
    i int,
    v varchar
);

INSERT INTO test (i, v) VALUES (1, 'first line');
INSERT INTO test (i, v) VALUES (2, 'second line');
INSERT INTO test (i, v) VALUES (3, 'third line');
INSERT INTO test (i, v) VALUES (4, 'immortal');

CREATE OR REPLACE FUNCTION test_munge() RETURNS SETOF test AS $$
    my $rv = spi_exec_query('select i, v from test;');
    my $status = $rv->{'status'};
    my $nrows = $rv->{'processed'};
    foreach my $rn (0 .. $nrows - 1) {
        my $row = $rv->{'rows'}[$rn];
        $row->{'i'} += 200 if defined($row->{'i'});
        $row->{'v'} =~ tr/A-Za-z/a-zA-Z/ if (defined($row->{'v'}));
        return_next($row);
    }
    return undef;
$$ LANGUAGE plperl;
SELECT * FROM test_munge();
```

```
spi_query(command)
spi_fetchrow(cursor)
spi_cursor_close(cursor)
```

`spi_query` and `spi_fetchrow` work together as a pair for row sets which might be large, or for cases where you wish to return rows as they arrive. `spi_fetchrow` works only with `spi_query`. The following example illustrates how you use them together:

```
CREATE TYPE foo_type AS (the_num INTEGER, the_text TEXT);
```
CREATE OR REPLACE FUNCTION lotsa_md5(INTEGER) RETURNS SETOF foo_type AS $$
use Digest::MD5 qw(md5_hex);
my $file = '/usr/share/dict/words';
my $t = localtime;
elog(NOTICE, "opening file $file at $t");
open my $fh, '<', $file or elog(ERROR, "cannot open $file for reading: $!");
my @words = <$fh>;
close $fh;
$t = localtime;
elog(NOTICE, "closed file $file at $t");
chomp(@words);
my $row;
my $sth = spi_query("SELECT * FROM generate_series(1,$_[0])
AS b(a)");
while (defined ($row = spi_fetchrow($sth))) {
  return_next({
    the_num => $row->{a},
    the_text => md5_hex($words[rand @words])
  });
}
return;
$$ LANGUAGE plperlu;

SELECT * from lotsa_md5(500);

Normally, spi_fetchrow should be repeated until it returns undef, indicating that there
are no more rows to read. The cursor returned by spi_query is automatically freed when
spi_fetchrow returns undef. If you do not wish to read all the rows, instead call spi_cur-
sor_close to free the cursor. Failure to do so will result in memory leaks.

spi_prepare(command, argument types)
spi_query_prepared(plan, arguments)
spi_exec_prepared(plan [, attributes], arguments)
spi_freeplan(plan)

spi_prepare, spi_query_prepared, spi_exec_prepared, and spi_freeplan
implement the same functionality but for prepared queries. spi_prepare accepts a query string
with numbered argument placeholders ($1, $2, etc) and a string list of argument types:

$plan = spi_prepare('SELECT * FROM test WHERE id > $1 AND name = $2',
                   'INTEGER',
                   'TEXT');

Once a query plan is prepared by a call to spi_prepare, the plan can be used instead of
the string query, either in spi_exec_prepared, where the result is the same as returned
by spi_exec_query, or in spi_query_prepared which returns a cursor exactly as
spi_query does, which can be later passed to spi_fetchrow. The optional second param-
ter to spi_exec_prepared is a hash reference of attributes; the only attribute currently sup-
ported is limit, which sets the maximum number of rows returned by a query.

The advantage of prepared queries is that it is possible to use one prepared plan for more than one
query execution. After the plan is not needed anymore, it can be freed with spi_freeplan:

CREATE OR REPLACE FUNCTION init() RETURNS VOID AS $$
$_SHARED{my_plan} = spi_prepare('SELECT (now() + $1)::date AS now', 'INTERVAL');

CREATE OR REPLACE FUNCTION add_time( INTERVAL ) RETURNS TEXT AS $$
    return spi_exec_prepared($$_SHARED{my_plan}, $$[0] )->{rows}->[0]->{now};
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION done() RETURNS VOID AS $$
    spi_freeplan( $$._SHARED{my_plan});
    undef $$._SHARED{my_plan};
$$ LANGUAGE plperl;

SELECT init();
SELECT add_time('1 day'), add_time('2 days'), add_time('3 days');
SELECT done();

    add_time | add_time | add_time
--------------+----------+----------
    2005-12-10 | 2005-12-11 | 2005-12-12

Note that the parameter subscript in spi_prepare is defined via $1, $2, $3, etc, so avoid declaring query strings in double quotes that might easily lead to hard-to-catch bugs.

Another example illustrates usage of an optional parameter in spi_exec_prepared:

CREATE TABLE hosts AS SELECT id, ('192.168.1.'||id)::inet AS address FROM generate_series(1,3) AS id;

CREATE OR REPLACE FUNCTION init_hosts_query() RETURNS VOID AS $$
    $$._SHARED{plan} = spi_prepare('SELECT * FROM hosts WHERE address << $1', 'inet');
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION query_hosts(inet) RETURNS SETOF hosts AS $$
    return spi_exec_prepared($$_SHARED{plan}, {limit => 2}, $$[0] )->{rows};
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION release_hosts_query() RETURNS VOID AS $$
    spi_freeplan($$_SHARED{plan});
    undef $$._SHARED{plan};
$$ LANGUAGE plperl;

SELECT init_hosts_query();
SELECT query_hosts('192.168.1.0/30');
SELECT release_hosts_query();

<table>
<thead>
<tr>
<th>query_hosts</th>
</tr>
</thead>
<tbody>
<tr>
<td>(1, 192.168.1.1)</td>
</tr>
<tr>
<td>(2, 192.168.1.2)</td>
</tr>
<tr>
<td>(2 rows)</td>
</tr>
</tbody>
</table>

spi_commit()
spi_rollback()

Commit or roll back the current transaction. This can only be called in a procedure or anonymous code block (DO command) called from the top level. (Note that it is not possible to run the SQL commands COMMIT or ROLLBACK via spi_exec_query or similar. It has to be done using these functions.) After a transaction is ended, a new transaction is automatically started, so there is no separate function for that.

Here is an example:

```sql
CREATE PROCEDURE transaction_test1()
LANGUAGE plperl
AS $$
    foreach my $i (0..9) {
        spi_exec_query("INSERT INTO test1 (a) VALUES ($i)");
        if ($i % 2 == 0) {
            spi_commit();
        } else {
            spi_rollback();
        }
    }
$$;
CALL transaction_test1();
```

### 45.3.2. Utility Functions in PL/Perl

**elog(level, msg)**

Emit a log or error message. Possible levels are DEBUG, LOG, INFO, NOTICE, WARNING, and ERROR. ERROR raises an error condition; if this is not trapped by the surrounding Perl code, the error propagates out to the calling query, causing the current transaction or subtransaction to be aborted. This is effectively the same as the Perl `die` command. The other levels only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the `log_min_messages` and `client_min_messages` configuration variables. See Chapter 19 for more information.

**quote_literal(string)**

Return the given string suitably quoted to be used as a string literal in an SQL statement string. Embedded single-quotes and backslashes are properly doubled. Note that `quote_literal` returns `undef` on `undef` input; if the argument might be `undef`, `quote_nullable` is often more suitable.

**quote_nullable(string)**

Return the given string suitably quoted to be used as a string literal in an SQL statement string; or, if the argument is `undef`, return the unquoted string "NULL". Embedded single-quotes and backslashes are properly doubled.
quote_ident(string)

Return the given string suitably quoted to be used as an identifier in an SQL statement string. Quotes are added only if necessary (i.e., if the string contains non-identifier characters or would be case-folded). Embedded quotes are properly doubled.

decode_bytea(string)

Return the unescaped binary data represented by the contents of the given string, which should be bytea encoded.

encode_bytea(string)

Return the bytea encoded form of the binary data contents of the given string.

encode_array_literal(array)

encode_array_literal(array, delimiter)

Returns the contents of the referenced array as a string in array literal format (see Section 8.15.2). Returns the argument value unaltered if it's not a reference to an array. The delimiter used between elements of the array literal defaults to "," if a delimiter is not specified or is undef.

encode_typed_literal(value, typename)

Converts a Perl variable to the value of the data type passed as a second argument and returns a string representation of this value. Correctly handles nested arrays and values of composite types.

encode_array_constructor(array)

Returns the contents of the referenced array as a string in array constructor format (see Section 4.2.12). Individual values are quoted using quote_nullable. Returns the argument value, quoted using quote_nullable, if it's not a reference to an array.

looks_like_number(string)

Returns a true value if the content of the given string looks like a number, according to Perl, returns false otherwise. Returns undef if the argument is undef. Leading and trailing space is ignored. Inf and Infinity are regarded as numbers.

is_array_ref(argument)

Returns a true value if the given argument may be treated as an array reference, that is, if ref of the argument is ARRAY or PostgreSQL::InServer::ARRAY. Returns false otherwise.

45.4. Global Values in PL/Perl

You can use the global hash %_SHARED to store data, including code references, between function calls for the lifetime of the current session.

Here is a simple example for shared data:

CREATE OR REPLACE FUNCTION set_var(name text, val text) RETURNS text AS $$
    if ($_SHARED{$_[0]} = $_[1]) {
        return 'ok';
    } else {
        return "cannot set shared variable $_[0] to $_[1]";
    }
$$ LANGUAGE plperl;
CREATE OR REPLACE FUNCTION get_var(name text) RETURNS text AS $$
return $_[0];
$$ LANGUAGE plperl;

SELECT set_var('sample', 'Hello, PL/Perl! How''s tricks?');
SELECT get_var('sample');

Here is a slightly more complicated example using a code reference:

CREATE OR REPLACE FUNCTION myfuncs() RETURNS void AS $$
$_SHARED{myquote} = sub {
    my $arg = shift;
    $arg =~ s/\['\\]/\$1/g;
    return "$arg";
}
$$ LANGUAGE plperl;

SELECT myfuncs(); /* initializes the function */

/* Set up a function that uses the quote function */

CREATE OR REPLACE FUNCTION use_quote(TEXT) RETURNS text AS $$
    my $text_to_quote = shift;
    my $qfunc = $_[0];
    return &$qfunc($text_to_quote);
$$ LANGUAGE plperl;

(You could have replaced the above with the one-liner return $_[0]; at the expense of readability.)

For security reasons, PL/Perl executes functions called by any one SQL role in a separate Perl interpreter for that role. This prevents accidental or malicious interference by one user with the behavior of another user's PL/Perl functions. Each such interpreter has its own value of the $_SHARED variable and other global state. Thus, two PL/Perl functions will share the same value of $_SHARED if and only if they are executed by the same SQL role. In an application wherein a single session executes code under multiple SQL roles (via SECURITY DEFINER functions, use of SET ROLE, etc) you may need to take explicit steps to ensure that PL/Perl functions can share data via $_SHARED. To do that, make sure that functions that should communicate are owned by the same user, and mark them SECURITY DEFINER. You must of course take care that such functions can't be used to do anything unintended.

45.5. Trusted and Untrusted PL/Perl

Normally, PL/Perl is installed as a “trusted” programming language named plperl. In this setup, certain Perl operations are disabled to preserve security. In general, the operations that are restricted are those that interact with the environment. This includes file handle operations, require, and use (for external modules). There is no way to access internals of the database server process or to gain OS-level access with the permissions of the server process, as a C function can do. Thus, any unprivileged database user can be permitted to use this language.

Here is an example of a function that will not work because file system operations are not allowed for security reasons:

CREATE FUNCTION badfunc() RETURNS integer AS $$
    my $tmpfile = "/tmp/badfile";
    open my $fh, '>', $tmpfile
$$ LANGUAGE plperl;
PL/Perl - Perl Procedural Language

```perl
or elog(ERROR, qq{could not open the file "$tmpfile": $!});
print $fh "Testing writing to a file\n";
close $fh or elog(ERROR, qq{could not close the file
"$tmpfile": $!});
return 1;
```

The creation of this function will fail as its use of a forbidden operation will be caught by the validator.

Sometimes it is desirable to write Perl functions that are not restricted. For example, one might want a Perl function that sends mail. To handle these cases, PL/Perl can also be installed as an “untrusted” language (usually called PL/PerlU). In this case the full Perl language is available. When installing the language, the language name `plperl` will select the untrusted PL/Perl variant.

The writer of a PL/PerlU function must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator. Note that the database system allows only database superusers to create functions in untrusted languages.

If the above function was created by a superuser using the language `plperl`, execution would succeed.

In the same way, anonymous code blocks written in Perl can use restricted operations if the language is specified as `plperl` rather than `plperl`, but the caller must be a superuser.

```
Note

While PL/Perl functions run in a separate Perl interpreter for each SQL role, all PL/PerlU functions executed in a given session run in a single Perl interpreter (which is not any of the ones used for PL/Perl functions). This allows PL/PerlU functions to share data freely, but no communication can occur between PL/Perl and PL/PerlU functions.

Note

Perl cannot support multiple interpreters within one process unless it was built with the appropriate flags, namely either `usemultiplicity` or `useithreads`. (`usemultiplicity` is preferred unless you actually need to use threads. For more details, see the `perlembed` man page.) If PL/Perl is used with a copy of Perl that was not built this way, then it is only possible to have one Perl interpreter per session, and so any one session can only execute either PL/PerlU functions, or PL/Perl functions that are all called by the same SQL role.
```

### 45.6. PL/Perl Triggers

PL/Perl can be used to write trigger functions. In a trigger function, the hash reference `$_TD` contains information about the current trigger event. `$_TD` is a global variable, which gets a separate local value for each invocation of the trigger. The fields of the `$_TD` hash reference are:

- `$_TD->{new}{foo}`
  - NEW value of column `foo`
- `$_TD->{old}{foo}`
  - OLD value of column `foo`
$_TD->{name}
   Name of the trigger being called

$_TD->{event}
   Trigger event: INSERT, UPDATE, DELETE, TRUNCATE, or UNKNOWN

$_TD->{when}
   When the trigger was called: BEFORE, AFTER, INSTEAD OF, or UNKNOWN

$_TD->{level}
   The trigger level: ROW, STATEMENT, or UNKNOWN

$_TD->{relid}
   OID of the table on which the trigger fired

$_TD->{table_name}
   Name of the table on which the trigger fired

$_TD->{relname}
   Name of the table on which the trigger fired. This has been deprecated, and could be removed in a future release. Please use $_TD->{table_name} instead.

$_TD->{table_schema}
   Name of the schema in which the table on which the trigger fired, is

$_TD->{argc}
   Number of arguments of the trigger function

@{$_TD->{args}}
   Arguments of the trigger function. Does not exist if $_TD->{argc} is 0.

Row-level triggers can return one of the following:

  return;
   Execute the operation

  "SKIP"
   Don't execute the operation

  "MODIFY"
   Indicates that the NEW row was modified by the trigger function

Here is an example of a trigger function, illustrating some of the above:

CREATE TABLE test (
   i int,
   v varchar
);

CREATE OR REPLACE FUNCTION valid_id() RETURNS trigger AS $$
if (($_TD->{new}{i} >= 100) || ($_TD->{new}{i} <= 0)) {
    return "SKIP";    # skip INSERT/UPDATE command
} elsif ($_TD->{new}{v} ne "immortal") {
    $_TD->{new}{v} .= "(modified by trigger)";
    return "MODIFY";  # modify row and execute INSERT/UPDATE command
} else {
    return;           # execute INSERT/UPDATE command
}
$$ LANGUAGE plperl;

CREATE TRIGGER test_valid_id_trig
BEFORE INSERT OR UPDATE ON test
FOR EACH ROW EXECUTE FUNCTION valid_id();

45.7. PL/Perl Event Triggers

PL/Perl can be used to write event trigger functions. In an event trigger function, the hash reference
$_TD contains information about the current trigger event. $_TD is a global variable, which gets a
separate local value for each invocation of the trigger. The fields of the $_TD hash reference are:

$_TD->{event}
The name of the event the trigger is fired for.

$_TD->{tag}
The command tag for which the trigger is fired.

The return value of the trigger function is ignored.

Here is an example of an event trigger function, illustrating some of the above:

CREATE OR REPLACE FUNCTION perlsnitch() RETURNS event_trigger AS $$
  elog(NOTICE, "perlsnitch: " . $_TD->{event} . " " . $_TD->{tag} . " ");
$$ LANGUAGE plperl;

CREATE EVENT TRIGGER perl_a_snitch
ON ddl_command_start
  EXECUTE FUNCTION perlsnitch();

45.8. PL/Perl Under the Hood

45.8.1. Configuration

This section lists configuration parameters that affect PL/Perl.

plperl.on_init(string)

Specifies Perl code to be executed when a Perl interpreter is first initialized, before it is specialized
for use by plperl or plperlu. The SPI functions are not available when this code is executed.
If the code fails with an error it will abort the initialization of the interpreter and propagate out to
the calling query, causing the current transaction or subtransaction to be aborted.

The Perl code is limited to a single string. Longer code can be placed into a module and loaded
by the on_init string. Examples:
plperl.on_init = 'require "plperlinit.pl"'
plperl.on_init = 'use lib "/my/app"; use MyApp::PgInit;'

Any modules loaded by plperl.on_init, either directly or indirectly, will be available for use by plperl. This may create a security risk. To see what modules have been loaded you can use:

DO 'elog(WARNING, join ", ", sort keys %INC)' LANGUAGE plperl;

Initialization will happen in the postmaster if the plperl library is included in shared_preload_libraries, in which case extra consideration should be given to the risk of destabilizing the postmaster. The principal reason for making use of this feature is that Perl modules loaded by plperl.on_init need be loaded only at postmaster start, and will be instantly available without loading overhead in individual database sessions. However, keep in mind that the overhead is avoided only for the first Perl interpreter used by a database session — either PL/PerlU, or PL/Perl for the first SQL role that calls a PL/Perl function. Any additional Perl interpreters created in a database session will have to execute plperl.on_init afresh. Also, on Windows there will be no savings whatsoever from preloading, since the Perl interpreter created in the postmaster process does not propagate to child processes.

This parameter can only be set in the postgresql.conf file or on the server command line.

plperl.on_plperl_init (string)
plperl.on_plperlu_init (string)

These parameters specify Perl code to be executed when a Perl interpreter is specialized for plperl or plperlu respectively. This will happen when a PL/Perl or PL/PerlU function is first executed in a database session, or when an additional interpreter has to be created because the other language is called or a PL/Perl function is called by a new SQL role. This follows any initialization done by plperl.on_init. The SPI functions are not available when this code is executed. The Perl code in plperl.on_plperl_init is executed after “locking down” the interpreter, and thus it can only perform trusted operations.

If the code fails with an error it will abort the initialization and propagate out to the calling query, causing the current transaction or subtransaction to be aborted. Any actions already done within Perl won't be undone; however, that interpreter won't be used again. If the language is used again the initialization will be attempted again within a fresh Perl interpreter.

Only superusers can change these settings. Although these settings can be changed within a session, such changes will not affect Perl interpreters that have already been used to execute functions.

plperl.use_strict (boolean)

When set true subsequent compilations of PL/Perl functions will have the strict pragma enabled. This parameter does not affect functions already compiled in the current session.

45.8.2. Limitations and Missing Features

The following features are currently missing from PL/Perl, but they would make welcome contributions.

• PL/Perl functions cannot call each other directly.

• SPI is not yet fully implemented.

• If you are fetching very large data sets using spi_exec_query, you should be aware that these will all go into memory. You can avoid this by using spi_query/spi_fetchrow as illustrated earlier.
A similar problem occurs if a set-returning function passes a large set of rows back to PostgreSQL via return. You can avoid this problem too by instead using return_next for each row returned, as shown previously.

- When a session ends normally, not due to a fatal error, any END blocks that have been defined are executed. Currently no other actions are performed. Specifically, file handles are not automatically flushed and objects are not automatically destroyed.
Chapter 46. PL/Python - Python Procedural Language

The PL/Python procedural language allows PostgreSQL functions to be written in the Python language. To install PL/Python in a particular database, use CREATE EXTENSION plpythonu (but see also Section 46.1).

Tip

If a language is installed into template1, all subsequently created databases will have the language installed automatically.

PL/Python is only available as an “untrusted” language, meaning it does not offer any way of restricting what users can do in it and is therefore named plpythonu. A trusted variant plpython might become available in the future if a secure execution mechanism is developed in Python. The writer of a function in untrusted PL/Python must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator. Only superusers can create functions in untrusted languages such as plpythonu.

Note

Users of source packages must specially enable the build of PL/Python during the installation process. (Refer to the installation instructions for more information.) Users of binary packages might find PL/Python in a separate subpackage.

46.1. Python 2 vs. Python 3

PL/Python supports both the Python 2 and Python 3 language variants. (The PostgreSQL installation instructions might contain more precise information about the exact supported minor versions of Python.) Because the Python 2 and Python 3 language variants are incompatible in some important aspects, the following naming and transitioning scheme is used by PL/Python to avoid mixing them:

• The PostgreSQL language named plpython2u implements PL/Python based on the Python 2 language variant.

• The PostgreSQL language named plpython3u implements PL/Python based on the Python 3 language variant.

• The language named plpythonu implements PL/Python based on the default Python language variant, which is currently Python 2. (This default is independent of what any local Python installations might consider to be their “default”, for example, what /usr/bin/python might be.) The default will probably be changed to Python 3 in a distant future release of PostgreSQL, depending on the progress of the migration to Python 3 in the Python community.

This scheme is analogous to the recommendations in PEP 394 regarding the naming and transitioning of the python command.

1 https://www.python.org
2 https://www.python.org/dev/peps/pep-394/
It depends on the build configuration or the installed packages whether PL/Python for Python 2 or Python 3 or both are available.

**Tip**

The built variant depends on which Python version was found during the installation or which version was explicitly set using the PYTHON environment variable; see Section 16.4. To make both variants of PL/Python available in one installation, the source tree has to be configured and built twice.

This results in the following usage and migration strategy:

- Existing users and users who are currently not interested in Python 3 use the language name `plpythonu` and don’t have to change anything for the foreseeable future. It is recommended to gradually “future-proof” the code via migration to Python 2.6/2.7 to simplify the eventual migration to Python 3.

- Users who want to dive into Python 3 can use the `plpython3u` language name, which will keep working forever by today’s standards. In the distant future, when Python 3 might become the default, they might like to remove the “3” for aesthetic reasons.

- Users who want to build a Python-3-only operating system environment, can change the contents of `pg_pltemplate` to make `plpythonu` be equivalent to `plpython3u`, keeping in mind that this would make their installation incompatible with most of the rest of the world.

See also the document What's New In Python 3.0 for more information about porting to Python 3.

It is not allowed to use PL/Python based on Python 2 and PL/Python based on Python 3 in the same session, because the symbols in the dynamic modules would clash, which could result in crashes of the PostgreSQL server process. There is a check that prevents mixing Python major versions in a session, which will abort the session if a mismatch is detected. It is possible, however, to use both PL/Python variants in the same database, from separate sessions.

### 46.2. PL/Python Functions

Functions in PL/Python are declared via the standard CREATE FUNCTION syntax:

```sql
CREATE FUNCTION funcname (argument-list)
    RETURNS return-type
AS $$
    # PL/Python function body
$$ LANGUAGE plpythonu;
```

The body of a function is simply a Python script. When the function is called, its arguments are passed as elements of the list `args`; named arguments are also passed as ordinary variables to the Python script. Use of named arguments is usually more readable. The result is returned from the Python code in the usual way, with `return` or `yield` (in case of a result-set statement). If you do not provide a return value, Python returns the default `None`. PL/Python translates Python’s `None` into the SQL null

---

3. [https://docs.python.org/3/whatsnew/3.0.html](https://docs.python.org/3/whatsnew/3.0.html)
value. In a procedure, the result from the Python code must be None (typically achieved by ending the procedure without a return statement or by using a return statement without argument); otherwise, an error will be raised.

For example, a function to return the greater of two integers can be defined as:

```python
CREATE FUNCTION pymax (a integer, b integer)
RETURNS integer
AS $$
if a > b:
    return a
return b
$$ LANGUAGE plpythonu;
```

The Python code that is given as the body of the function definition is transformed into a Python function. For example, the above results in:

```python
def __plpython_procedure_pymax_23456():
    if a > b:
        return a
    return b
```
assuming that 23456 is the OID assigned to the function by PostgreSQL.

The arguments are set as global variables. Because of the scoping rules of Python, this has the subtle consequence that an argument variable cannot be reassigned inside the function to the value of an expression that involves the variable name itself, unless the variable is redeclared as global in the block. For example, the following won’t work:

```python
CREATE FUNCTION pystrip(x text)
RETURNS text
AS $$
x = x.strip()  # error
return x
$$ LANGUAGE plpythonu;
```
because assigning to x makes x a local variable for the entire block, and so the x on the right-hand side of the assignment refers to a not-yet-assigned local variable x, not the PL/Python function parameter. Using the global statement, this can be made to work:

```python
CREATE FUNCTION pystrip(x text)
RETURNS text
AS $$
global x
x = x.strip()  # ok now
return x
$$ LANGUAGE plpythonu;
```

But it is advisable not to rely on this implementation detail of PL/Python. It is better to treat the function parameters as read-only.

### 46.3. Data Values

Generally speaking, the aim of PL/Python is to provide a “natural” mapping between the PostgreSQL and the Python worlds. This informs the data mapping rules described below.
46.3.1. Data Type Mapping

When a PL/Python function is called, its arguments are converted from their PostgreSQL data type to a corresponding Python type:

- PostgreSQL `boolean` is converted to Python `bool`.
- PostgreSQL `smallint` and `int` are converted to Python `int`. PostgreSQL `bigint` and `oid` are converted to `long` in Python 2 and to `int` in Python 3.
- PostgreSQL `real` and `double` are converted to Python `float`.
- PostgreSQL `numeric` is converted to Python `Decimal`. This type is imported from the `cdecimal` package if that is available. Otherwise, `decimal.Decimal` from the standard library will be used. `cdecimal` is significantly faster than `decimal`. In Python 3.3 and up, however, `cdecimal` has been integrated into the standard library under the name `decimal`, so there is no longer any difference.
- PostgreSQL `bytea` is converted to Python `str` in Python 2 and to `bytes` in Python 3. In Python 2, the string should be treated as a byte sequence without any character encoding.
- All other data types, including the PostgreSQL character string types, are converted to a Python `str`. In Python 2, this string will be in the PostgreSQL server encoding; in Python 3, it will be a Unicode string like all strings.
- For nonscallar data types, see below.

When a PL/Python function returns, its return value is converted to the function's declared PostgreSQL return data type as follows:

- When the PostgreSQL return type is `boolean`, the return value will be evaluated for truth according to the Python rules. That is, 0 and empty string are false, but notably `'f'` is true.
- When the PostgreSQL return type is `bytea`, the return value will be converted to a string (Python 2) or `bytes` (Python 3) using the respective Python built-ins, with the result being converted to `bytea`.
- For all other PostgreSQL return types, the return value is converted to a string using the Python built-in `str`, and the result is passed to the input function of the PostgreSQL data type. (If the Python value is a `float`, it is converted using the `repr` built-in instead of `str`, to avoid loss of precision.)

Strings in Python 2 are required to be in the PostgreSQL server encoding when they are passed to PostgreSQL. Strings that are not valid in the current server encoding will raise an error, but not all encoding mismatches can be detected, so garbage data can still result when this is not done correctly. Unicode strings are converted to the correct encoding automatically, so it can be safer and more convenient to use those. In Python 3, all strings are Unicode strings.

- For nonscallar data types, see below.

Note that logical mismatches between the declared PostgreSQL return type and the Python data type of the actual return object are not flagged; the value will be converted in any case.

46.3.2. Null, None

If an SQL null value is passed to a function, the argument value will appear as `None` in Python. For example, the function definition of `pymax` shown in Section 46.2 will return the wrong answer for null inputs. We could add `STRICT` to the function definition to make PostgreSQL do something more reasonable: if a null value is passed, the function will not be called at all, but will just return a null result automatically. Alternatively, we could check for null inputs in the function body:
CREATE FUNCTION pymax (a integer, b integer)
    RETURNS integer
AS $$
if (a is None) or (b is None):
    return None
if a > b:
    return a
return b
$$ LANGUAGE plpythonu;

As shown above, to return an SQL null value from a PL/Python function, return the value None. This can be done whether the function is strict or not.

### 46.3.3. Arrays, Lists

SQL array values are passed into PL/Python as a Python list. To return an SQL array value out of a PL/Python function, return a Python list:

```
CREATE FUNCTION return_arr()
    RETURNS int[]
AS $$
return [1, 2, 3, 4, 5]
$$ LANGUAGE plpythonu;

SELECT return_arr();
```

```
return_arr
--------
{1,2,3,4,5}
(1 row)
```

Multidimensional arrays are passed into PL/Python as nested Python lists. A 2-dimensional array is a list of lists, for example. When returning a multi-dimensional SQL array out of a PL/Python function, the inner lists at each level must all be of the same size. For example:

```
CREATE FUNCTION test_type_conversion_array_int4(x int4[]) RETURNS int4[] AS $$
plpy.info(x, type(x))
return x
$$ LANGUAGE plpythonu;

SELECT * FROM test_type_conversion_array_int4(ARRAY[[1,2,3], [4,5,6]]);
```

```
INFO:  ([1, 2, 3], [4, 5, 6], <type 'list'>)
test_type_conversion_array_int4
------------------------
[1,2,3],[4,5,6]
(1 row)
```

Other Python sequences, like tuples, are also accepted for backwards-compatibility with PostgreSQL versions 9.6 and below, when multi-dimensional arrays were not supported. However, they are always treated as one-dimensional arrays, because they are ambiguous with composite types. For the same reason, when a composite type is used in a multi-dimensional array, it must be represented by a tuple, rather than a list.

Note that in Python, strings are sequences, which can have undesirable effects that might be familiar to Python programmers:
CREATE FUNCTION return_str_arr()
    RETURNS varchar[]
AS $$
return "hello"
$$ LANGUAGE plpythonu;

SELECT return_str_arr();

return_str_arr
----------------
{h,e,l,l,o}
(1 row)

46.3.4. Composite Types

Composite-type arguments are passed to the function as Python mappings. The element names of the mapping are the attribute names of the composite type. If an attribute in the passed row has the null value, it has the value None in the mapping. Here is an example:

CREATE TABLE employee (  
    name text,  
    salary integer,  
    age integer  
);

CREATE FUNCTION overpaid (e employee)
    RETURNS boolean
AS $$
    if e["salary"] > 200000:
        return True
    if (e["age"] < 30) and (e["salary"] > 100000):
        return True
    return False
$$ LANGUAGE plpythonu;

There are multiple ways to return row or composite types from a Python function. The following examples assume we have:

CREATE TYPE named_value AS (  
    name text,  
    value integer  
);

A composite result can be returned as a:

Sequence type (a tuple or list, but not a set because it is not indexable)

Returned sequence objects must have the same number of items as the composite result type has fields. The item with index 0 is assigned to the first field of the composite type, 1 to the second and so on. For example:

CREATE FUNCTION make_pair (name text, value integer)
    RETURNS named_value
AS $$
    return ( name, value )
    # or alternatively, as tuple: return [ name, value ]
$$\text{LANGUAGE plpythonu;}$$

To return a SQL null for any column, insert \texttt{None} at the corresponding position.

When an array of composite types is returned, it cannot be returned as a list, because it is ambiguous whether the Python list represents a composite type, or another array dimension.

**Mapping (dictionary)**

The value for each result type column is retrieved from the mapping with the column name as key. Example:

```sql
CREATE FUNCTION make_pair (name text, value integer) RETURNS named_value
AS $$
return { "name": name, "value": value }
$$ LANGUAGE plpythonu;
```

Any extra dictionary key/value pairs are ignored. Missing keys are treated as errors. To return a SQL null value for any column, insert \texttt{None} with the corresponding column name as the key.

**Object (any object providing method \texttt{__getattr__})**

This works the same as a mapping. Example:

```sql
CREATE FUNCTION make_pair (name text, value integer) RETURNS named_value
AS $$
class named_value:
    def __init__ (self, n, v):
        self.name = n
        self.value = v
    return named_value(name, value)

# or simply
class nv: pass
nv.name = name
nv.value = value
return nv
$$ LANGUAGE plpythonu;
```

Functions with \texttt{OUT} parameters are also supported. For example:

```sql
CREATE FUNCTION multiout_simple(OUT i integer, OUT j integer) AS $$
return (1, 2)
$$ LANGUAGE plpythonu;
```

```sql
SELECT * FROM multiout_simple();
```

Output parameters of procedures are passed back the same way. For example:

```sql
CREATE PROCEDURE python_triple(INOUT a integer, INOUT b integer) AS $$
return (a * 3, b * 3)
$$ LANGUAGE plpythonu;
```

```sql
CALL python_triple(5, 10);
```
46.3.5. Set-returning Functions

A PL/Python function can also return sets of scalar or composite types. There are several ways to achieve this because the returned object is internally turned into an iterator. The following examples assume we have composite type:

```sql
CREATE TYPE greeting AS (  
    how text,  
    who text
);
```

A set result can be returned from a:

Sequence type (tuple, list, set)

```sql
CREATE FUNCTION greet (how text)  
RETURNS SETOF greeting  
AS $\$$  
    # return tuple containing lists as composite types  
    # all other combinations work also  
    return ( [ how, "World" ], [ how, "PostgreSQL" ], [ how, "PL/Python" ] )  
$\$$ LANGUAGE plpythonu;
```

Iterator (any object providing `__iter__` and `next` methods)

```sql
CREATE FUNCTION greet (how text)  
RETURNS SETOF greeting  
AS $\$$  
    class producer:  
        def __init__ (self, how, who):  
            self.how = how  
            self.who = who  
            self.ndx = -1

        def __iter__ (self):  
            return self

        def next (self):  
            self.ndx += 1  
            if self.ndx == len(self.who):  
                raise StopIteration  
            return ( self.how, self.who[self.ndx] )

    return producer(how, [ "World", "PostgreSQL", "PL/Python" ])
$\$$ LANGUAGE plpythonu;
```

Generator (`yield`)

```sql
CREATE FUNCTION greet (how text)  
RETURNS SETOF greeting  
AS $\$$  
    for who in [ "World", "PostgreSQL", "PL/Python" ]:  
        yield ( how, who )  
$\$$ LANGUAGE plpythonu;
```
Set-returning functions with **OUT** parameters (using **RETURNS SETOF record**) are also supported. For example:

```plpython
CREATE FUNCTION multiout_simple_setof(n integer, OUT integer, OUT integer) RETURNS SETOF record AS $$
return [(1, 2)] * n
$$ LANGUAGE plpythonu;
```

```sql
SELECT * FROM multiout_simple_setof(3);
```

### 46.4. Sharing Data

The global dictionary **SD** is available to store private data between repeated calls to the same function. The global dictionary **GD** is public data, that is available to all Python functions within a session; use with care.

Each function gets its own execution environment in the Python interpreter, so that global data and function arguments from `myfunc` are not available to `myfunc2`. The exception is the data in the **GD** dictionary, as mentioned above.

### 46.5. Anonymous Code Blocks

PL/Python also supports anonymous code blocks called with the **DO** statement:

```plpython
DO $$
    # PL/Python code
$$ LANGUAGE plpythonu;
```

An anonymous code block receives no arguments, and whatever value it might return is discarded. Otherwise it behaves just like a function.

### 46.6. Trigger Functions

When a function is used as a trigger, the dictionary **TD** contains trigger-related values:

**TD["event"]**

contains the event as a string: **INSERT**, **UPDATE**, **DELETE**, or **TRUNCATE**.

**TD["when"]**

contains one of **BEFORE**, **AFTER**, or **INSTEAD OF**.

**TD["level"]**

contains **ROW** or **STATEMENT**.

**TD["new"]**

**TD["old"]**

For a row-level trigger, one or both of these fields contain the respective trigger rows, depending on the trigger event.

**TD["name"]**

contains the trigger name.
contains the name of the table on which the trigger occurred.

contains the schema of the table on which the trigger occurred.

contains the OID of the table on which the trigger occurred.

If the CREATE TRIGGER command included arguments, they are available in TD["args"][0] to TD["args"][n-1].

If TD["when"] is BEFORE or INSTEAD OF and TD["level"] is ROW, you can return None or "OK" from the Python function to indicate the row is unmodified, "SKIP" to abort the event, or if TD["event"] is INSERT or UPDATE you can return "MODIFY" to indicate you've modified the new row. Otherwise the return value is ignored.

### 46.7. Database Access

The PL/Python language module automatically imports a Python module called `plpy`. The functions and constants in this module are available to you in the Python code as `plpy.foo`.

### 46.7.1. Database Access Functions

The `plpy` module provides several functions to execute database commands:

```
plpy.execute(query [, max-rows])
```

Calling `plpy.execute` with a query string and an optional row limit argument causes that query to be run and the result to be returned in a result object.

The result object emulates a list or dictionary object. The result object can be accessed by row number and column name. For example:

```
rv = plpy.execute("SELECT * FROM my_table", 5)
```

returns up to 5 rows from `my_table`. If `my_table` has a column `my_column`, it would be accessed as:

```
foo = rv[i]["my_column"]
```

The number of rows returned can be obtained using the built-in `len` function.

The result object has these additional methods:

```
nrows()
```

Returns the number of rows processed by the command. Note that this is not necessarily the same as the number of rows returned. For example, an `UPDATE` command will set this value but won't return any rows (unless `RETURNING` is used).

```
status()
```

The SPI_execute() return value.
colnames()
coltypes()
coltypmods()

Return a list of column names, list of column type OIDs, and list of type-specific type modifiers for the columns, respectively.

These methods raise an exception when called on a result object from a command that did not produce a result set, e.g., UPDATE without RETURNING, or DROP TABLE. But it is OK to use these methods on a result set containing zero rows.

__str__()

The standard __str__ method is defined so that it is possible for example to debug query execution results using plpy.debug(rv).

The result object can be modified.

Note that calling plpy.execute will cause the entire result set to be read into memory. Only use that function when you are sure that the result set will be relatively small. If you don't want to risk excessive memory usage when fetching large results, use plpy.cursor rather than plpy.execute.

plpy.prepare(query [, argtypes])
plpy.execute(plan [, arguments [, max-rows]])

plpy.prepare prepares the execution plan for a query. It is called with a query string and a list of parameter types, if you have parameter references in the query. For example:

plan = plpy.prepare("SELECT last_name FROM my_users WHERE first_name = $1", ["text"])
text is the type of the variable you will be passing for $1. The second argument is optional if you don't want to pass any parameters to the query.

After preparing a statement, you use a variant of the function plpy.execute to run it:

rv = plpy.execute(plan, ["name"], 5)

Pass the plan as the first argument (instead of the query string), and a list of values to substitute into the query as the second argument. The second argument is optional if the query does not expect any parameters. The third argument is the optional row limit as before.

Alternatively, you can call the execute method on the plan object:

rv = plan.execute(["name"], 5)

Query parameters and result row fields are converted between PostgreSQL and Python data types as described in Section 46.3.

When you prepare a plan using the PL/Python module it is automatically saved. Read the SPI documentation (Chapter 47) for a description of what this means. In order to make effective use of this across function calls one needs to use one of the persistent storage dictionaries SD or GD (see Section 46.4). For example:

CREATE FUNCTION usesavedplan() RETURNS trigger AS $$

1261
if "plan" in SD:
    plan = SD["plan"]
else:
    plan = plpy.prepare("SELECT 1")
    SD["plan"] = plan
  # rest of function
$$ LANGUAGE plpythonu;

plpy.cursor(query)
plpy.cursor(plan [, arguments])

The `plpy.cursor` function accepts the same arguments as `plpy.execute` (except for the row limit) and returns a cursor object, which allows you to process large result sets in smaller chunks. As with `plpy.execute`, either a query string or a plan object along with a list of arguments can be used, or the cursor function can be called as a method of the plan object.

The cursor object provides a `fetch` method that accepts an integer parameter and returns a result object. Each time you call `fetch`, the returned object will contain the next batch of rows, never larger than the parameter value. Once all rows are exhausted, `fetch` starts returning an empty result object. Cursor objects also provide an iterator interface, yielding one row at a time until all rows are exhausted. Data fetched that way is not returned as result objects, but rather as dictionaries, each dictionary corresponding to a single result row.

An example of two ways of processing data from a large table is:

```sql
CREATE FUNCTION count_odd_iterator() RETURNS integer AS $$
odd = 0
for row in plpy.cursor("select num from largetable"): 
  if row['num'] % 2:
    odd += 1
return odd
$$ LANGUAGE plpythonu;

CREATE FUNCTION count_odd_fetch(batch_size integer) RETURNS integer AS $$
odd = 0
cursor = plpy.cursor("select num from largetable")
while True:
    rows = cursor.fetch(batch_size)
    if not rows:
        break
    for row in rows:
        if row['num'] % 2:
            odd += 1
return odd
$$ LANGUAGE plpythonu;

CREATE FUNCTION count_odd_prepared() RETURNS integer AS $$
odd = 0
plan = plpy.prepare("select num from largetable where num % $1 <> 0", ["integer"])
rows = list(plpy.cursor(plan, [2]))  # or: = 
    list(plan.cursor([2]))
return len(rows)
$$ LANGUAGE plpythonu;
```

4 https://docs.python.org/library/stdtypes.html#iterator-types
Cursors are automatically disposed of. But if you want to explicitly release all resources held by a cursor, use the close method. Once closed, a cursor cannot be fetched from anymore.

**Tip**

Do not confuse objects created by `plpy.cursor` with DB-API cursors as defined by the Python Database API specification. They don't have anything in common except for the name.

### 46.7.2. Trapping Errors

Functions accessing the database might encounter errors, which will cause them to abort and raise an exception. Both `plpy.execute` and `plpy.prepare` can raise an instance of a subclass of `plpy.SPIError`, which by default will terminate the function. This error can be handled just like any other Python exception, by using the `try/except` construct. For example:

```plaintext
CREATE FUNCTION try_adding_joe() RETURNS text AS $$
    try:
        plpy.execute("INSERT INTO users(username) VALUES ('joe')")
    except plpy.SPIError:
        return "something went wrong"
    else:
        return "Joe added"
$$ LANGUAGE plpythonu;
```

The actual class of the exception being raised corresponds to the specific condition that caused the error. Refer to Table A.1 for a list of possible conditions. The module `plpy.spiexceptions` defines an exception class for each PostgreSQL condition, deriving their names from the condition name. For instance, `division_by_zero` becomes `DivisionByZero`, `unique_violation` becomes `UniqueViolation`, `fdw_error` becomes `FdwError`, and so on. Each of these exception classes inherits from `SPIError`. This separation makes it easier to handle specific errors, for instance:

```plaintext
CREATE FUNCTION insert_fraction(numerator int, denominator int)
    RETURNS text AS $$
    from plpy import spiexceptions
    try:
        plan = plpy.prepare("INSERT INTO fractions (frac) VALUES ($1 / $2)", ["int", "int"])
        plpy.execute(plan, [numerator, denominator])
    except spiexceptions.DivisionByZero:
        return "denominator cannot equal zero"
    except spiexceptions.UniqueViolation:
        return "already have that fraction"
    except plpy.SPIError, e:
        return "other error, SQLSTATE %s" % e.sqlstate
    else:
        return "fraction inserted"
$$ LANGUAGE plpythonu;
```

Note that because all exceptions from the `plpy.spiexceptions` module inherit from `SPIError`, an `except` clause handling it will catch any database access error.

---

5 https://www.python.org/dev/peps/pep-0249/
As an alternative way of handling different error conditions, you can catch the \texttt{SPIError} exception and determine the specific error condition inside the \texttt{except} block by looking at the \texttt{sqlstate} attribute of the exception object. This attribute is a string value containing the “SQLSTATE” error code. This approach provides approximately the same functionality.

### 46.8. Explicit Subtransactions

Recovering from errors caused by database access as described in Section 46.7.2 can lead to an undesirable situation where some operations succeed before one of them fails, and after recovering from that error the data is left in an inconsistent state. PL/Python offers a solution to this problem in the form of explicit subtransactions.

#### 46.8.1. Subtransaction Context Managers

Consider a function that implements a transfer between two accounts:

```python
CREATE FUNCTION transfer_funds() RETURNS void AS $$
try:
    plpy.execute("UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'")
    plpy.execute("UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'")
except plpy.SPIError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"
plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"]) plpy.execute(plan, [result])$$ LANGUAGE plpythonu;
```

If the second \texttt{UPDATE} statement results in an exception being raised, this function will report the error, but the result of the first \texttt{UPDATE} will nevertheless be committed. In other words, the funds will be withdrawn from Joe's account, but will not be transferred to Mary's account.

To avoid such issues, you can wrap your \texttt{plpy.execute} calls in an explicit subtransaction. The \texttt{plpy} module provides a helper object to manage explicit subtransactions that gets created with the \texttt{plpy.subtransaction()} function. Objects created by this function implement the context manager interface\(^6\). Using explicit subtransactions we can rewrite our function as:

```python
CREATE FUNCTION transfer_funds2() RETURNS void AS $$
try:
    with plpy.subtransaction():
        plpy.execute("UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'")
        plpy.execute("UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'")
except plpy.SPIError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"
plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"]) plpy.execute(plan, [result])$$ LANGUAGE plpythonu;
```

\(^6\) https://docs.python.org/library/stdtypes.html#context-manager-types
$$ LANGUAGE plpythonu;

Note that the use of try/catch is still required. Otherwise the exception would propagate to the top of the Python stack and would cause the whole function to abort with a PostgreSQL error, so that the operations table would not have any row inserted into it. The subtransaction context manager does not trap errors, it only assures that all database operations executed inside its scope will be atomically committed or rolled back. A rollback of the subtransaction block occurs on any kind of exception exit, not only ones caused by errors originating from database access. A regular Python exception raised inside an explicit subtransaction block would also cause the subtransaction to be rolled back.

46.8.2. Older Python Versions

Context managers syntax using the with keyword is available by default in Python 2.6. If using PL/Python with an older Python version, it is still possible to use explicit subtransactions, although not as transparently. You can call the subtransaction manager’s __enter__ and __exit__ functions using the enter and exit convenience aliases. The example function that transfers funds could be written as:

```
CREATE FUNCTION transfer_funds_old() RETURNS void AS $$
try:
    subxact = plpy.subtransaction()
    subxact.enter()
    try:
        plpy.execute("UPDATE accounts SET balance = balance - 100
WHERE account_name = 'joe'")
        plpy.execute("UPDATE accounts SET balance = balance + 100
WHERE account_name = 'mary'")
    except:
        import sys
        subxact.exit(*sys.exc_info())
        raise
    else:
        subxact.exit(None, None, None)
except plpy.SPIError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"
plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)",
                     ["text"])
plpy.execute(plan, [result])
$$ LANGUAGE plpythonu;
```

Note

Although context managers were implemented in Python 2.5, to use the with syntax in that version you need to use a future statement\(^7\). Because of implementation details, however, you cannot use future statements in PL/Python functions.

46.9. Transaction Management

In a procedure called from the top level or an anonymous code block (DO command) called from the top level it is possible to control transactions. To commit the current transaction, call plpy.commit().

\(^7\) [https://docs.python.org/release/2.5/ref/future.html](https://docs.python.org/release/2.5/ref/future.html)
To roll back the current transaction, call `plpy.rollback()`.
(Note that it is not possible to run the SQL commands `COMMIT` or `ROLLBACK` via `plpy.execute` or similar. It has to be done using these functions.) After a transaction is ended, a new transaction is automatically started, so there is no separate function for that.

Here is an example:

```sql
CREATE PROCEDURE transaction_test1()
LANGUAGE plpythonu
AS $$
for i in range(0, 10):
    plpy.execute("INSERT INTO test1 (a) VALUES (%d)" % i)
    if i % 2 == 0:
        plpy.commit()
    else:
        plpy.rollback()
$$;

CALL transaction_test1();
```

Transactions cannot be ended when an explicit subtransaction is active.

### 46.10. Utility Functions

The `plpy` module also provides the functions:

- `plpy.debug(msg, **kwargs)`
- `plpy.log(msg, **kwargs)`
- `plpy.info(msg, **kwargs)`
- `plpy.notice(msg, **kwargs)`
- `plpy.warning(msg, **kwargs)`
- `plpy.error(msg, **kwargs)`
- `plpy.fatal(msg, **kwargs)`

`plpy.error` and `plpy.fatal` actually raise a Python exception which, if uncaught, propagates out to the calling query, causing the current transaction or subtransaction to be aborted.

`raise plpy.Error(msg)` and `raise plpy.Fatal(msg)` are equivalent to calling `plpy.error(msg)` and `plpy.fatal(msg)`, respectively but the `raise` form does not allow passing keyword arguments. The other functions only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the `log_min_messages` and `client_min_messages` configuration variables. See Chapter 19 for more information.

The `msg` argument is given as a positional argument. For backward compatibility, more than one positional argument can be given. In that case, the string representation of the tuple of positional arguments becomes the message reported to the client.

The following keyword-only arguments are accepted:

- `detail`
- `hint`
- `sqlstate`
- `schema_name`
- `table_name`
- `column_name`
- `datatype_name`
- `constraint_name`
The string representation of the objects passed as keyword-only arguments is used to enrich the messages reported to the client. For example:

```plpython
CREATE FUNCTION raise_custom_exception() RETURNS void AS $$
plpy.error("custom exception message",
    detail="some info about exception",
    hint="hint for users")
$$ LANGUAGE plpythonu;

=# SELECT raise_custom_exception();
ERROR:  plpy.Error: custom exception message
DETAIL:  some info about exception
HINT:  hint for users
CONTEXT:  Traceback (most recent call last):
    PL/Python function "raise_custom_exception", line 4, in <module>
        hint="hint for users")
PL/Python function "raise_custom_exception"
```

Another set of utility functions are `plpy.quote_literal(string)`, `plpy.quote_nullable(string)` and `plpy.quote_ident(string)`. They are equivalent to the built-in quoting functions described in Section 9.4. They are useful when constructing ad-hoc queries. A PL/Python equivalent of dynamic SQL from Example 43.1 would be:

```plpython
plpy.execute("UPDATE tbl SET %s = %s WHERE key = %s" % (
    plpy.quote_ident(colname),
    plpy.quote_nullable(newvalue),
    plpy.quote_literal(keyvalue)))
```

### 46.11. Environment Variables

Some of the environment variables that are accepted by the Python interpreter can also be used to affect PL/Python behavior. They would need to be set in the environment of the main PostgreSQL server process, for example in a start script. The available environment variables depend on the version of Python; see the Python documentation for details. At the time of this writing, the following environment variables have an affect on PL/Python, assuming an adequate Python version:

- PYTHONHOME
- PYTHONPATH
- PYTHONY2K
- PYTHONOPTIMIZE
- PYTHONDEBUG
- PYTHONVERBOSE
- PYTHONCASEOK
- PYTHONDONTWRITEBYTECODE
- PYTHONIOENCODING
- PYTHONUSERBASE
- PYTHONHASHSEED
(It appears to be a Python implementation detail beyond the control of PL/Python that some of the environment variables listed on the python man page are only effective in a command-line interpreter and not an embedded Python interpreter.)
Chapter 47. Server Programming Interface

The Server Programming Interface (SPI) gives writers of user-defined C functions the ability to run SQL commands inside their functions. SPI is a set of interface functions to simplify access to the parser, planner, and executor. SPI also does some memory management.

Note

The available procedural languages provide various means to execute SQL commands from functions. Most of these facilities are based on SPI, so this documentation might be of use for users of those languages as well.

Note that if a command invoked via SPI fails, then control will not be returned to your C function. Rather, the transaction or subtransaction in which your C function executes will be rolled back. (This might seem surprising given that the SPI functions mostly have documented error-return conventions. Those conventions only apply for errors detected within the SPI functions themselves, however.) It is possible to recover control after an error by establishing your own subtransaction surrounding SPI calls that might fail.

SPI functions return a nonnegative result on success (either via a returned integer value or in the global variable SPI_result, as described below). On error, a negative result or NULL will be returned.

Source code files that use SPI must include the header file executor/spi.h.

47.1. Interface Functions
SPI_connect

SPI_connect, SPI_connect_ext — connect a C function to the SPI manager

Synopsis

```c
int SPI_connect(void)

int SPI_connect_ext(int options)
```

Description

SPI_connect opens a connection from a C function invocation to the SPI manager. You must call this function if you want to execute commands through SPI. Some utility SPI functions can be called from unconnected C functions.

SPI_connect_ext does the same but has an argument that allows passing option flags. Currently, the following option values are available:

- **SPI_OPT_NONATOMIC**

  Sets the SPI connection to be *nonatomic*, which means that transaction control calls SPI_commit, SPI_rollback, and SPI_start_transaction are allowed. Otherwise, calling these functions will result in an immediate error.

SPI_connect() is equivalent to SPI_connect_ext(0).

Return Value

- **SPI_OK_CONNECT**
  
  on success

- **SPI_ERROR_CONNECT**
  
  on error
SPI_finish

SPI_finish — disconnect a C function from the SPI manager

Synopsis

int SPI_finish(void)

Description

SPI_finish closes an existing connection to the SPI manager. You must call this function after completing the SPI operations needed during your C function's current invocation. You do not need to worry about making this happen, however, if you abort the transaction via \texttt{elog(\texttt{ERROR})}. In that case SPI will clean itself up automatically.

Return Value

SPI_OK_FINISH

if properly disconnected

SPI_ERROR_UNCONNECTED

if called from an unconnected C function
SPI_execute

SPI_execute — execute a command

Synopsis

int SPI_execute(const char * command, bool read_only, long count)

Description

SPI_execute executes the specified SQL command for count rows. If read_only is true, the command must be read-only, and execution overhead is somewhat reduced.

This function can only be called from a connected C function.

If count is zero then the command is executed for all rows that it applies to. If count is greater than zero, then no more than count rows will be retrieved; execution stops when the count is reached, much like adding a LIMIT clause to the query. For example,

SPI_execute("SELECT * FROM foo", true, 5);
will retrieve at most 5 rows from the table. Note that such a limit is only effective when the command actually returns rows. For example,

SPI_execute("INSERT INTO foo SELECT * FROM bar", false, 5);
inserts all rows from bar, ignoring the count parameter. However, with

SPI_execute("INSERT INTO foo SELECT * FROM bar RETURNING *", false, 5);
at most 5 rows would be inserted, since execution would stop after the fifth RETURNING result row is retrieved.

You can pass multiple commands in one string: SPI_execute returns the result for the command executed last. The count limit applies to each command separately (even though only the last result will actually be returned). The limit is not applied to any hidden commands generated by rules.

When read_only is false, SPI_execute increments the command counter and computes a new snapshot before executing each command in the string. The snapshot does not actually change if the current transaction isolation level is SERIALIZABLE or REPEATABLE READ, but in READ COMMITTED mode the snapshot update allows each command to see the results of newly committed transactions from other sessions. This is essential for consistent behavior when the commands are modifying the database.

When read_only is true, SPI_execute does not update either the snapshot or the command counter, and it allows only plain SELECT commands to appear in the command string. The commands are executed using the snapshot previously established for the surrounding query. This execution mode is somewhat faster than the read/write mode due to eliminating per-command overhead. It also allows genuinely stable functions to be built: since successive executions will all use the same snapshot, there will be no change in the results.

It is generally unwise to mix read-only and read-write commands within a single function using SPI; that could result in very confusing behavior, since the read-only queries would not see the results of any database updates done by the read-write queries.
The actual number of rows for which the (last) command was executed is returned in the global variable \texttt{SPI\_processed}. If the return value of the function is \texttt{SPI\_OK\_SELECT}, \texttt{SPI\_OK\_INSERT\_RETURNING}, \texttt{SPI\_OK\_DELETE\_RETURNING}, or \texttt{SPI\_OK\_UPDATE\_RETURNING}, then you can use the global pointer \texttt{SPI\_tuptable} to access the result rows. Some utility commands (such as \texttt{EXPLAIN}) also return row sets, and \texttt{SPI\_tuptable} will contain the result in these cases too. Some utility commands (\texttt{COPY}, \texttt{CREATE TABLE AS}) don't return a row set, so \texttt{SPI\_tuptable} is NULL, but they still return the number of rows processed in \texttt{SPI\_processed}.

The structure \texttt{SPITupleTable} is defined thus:

```c
typedef struct
{
    MemoryContext tuptabcxt;    /* memory context of result table */
    uint64      allocated;       /* number of allocated vals */
    uint64      free;            /* number of free vals */
    TupleDesc   tupdesc;         /* row descriptor */
    HeapTuple  *vals;            /* rows */
} SPITupleTable;
```

\texttt{vals} is an array of pointers to rows. (The number of valid entries is given by \texttt{SPI\_processed}.) \texttt{tupdesc} is a row descriptor which you can pass to SPI functions dealing with rows. \texttt{tuptabcxt}, \texttt{allocated}, and \texttt{free} are internal fields not intended for use by SPI callers.

\texttt{SPI\_finish} frees all \texttt{SPITupleTable}s allocated during the current C function. You can free a particular result table earlier, if you are done with it, by calling \texttt{SPI\_freetuptable}.

**Arguments**

- \texttt{const char * command}
  - string containing command to execute
- \texttt{bool read_only}
  - true for read-only execution
- \texttt{long count}
  - maximum number of rows to return, or 0 for no limit

**Return Value**

If the execution of the command was successful then one of the following (nonnegative) values will be returned:

- \texttt{SPI\_OK\_SELECT}
  - if a \texttt{SELECT} (but not \texttt{SELECT INTO}) was executed
- \texttt{SPI\_OK\_SELINTO}
  - if a \texttt{SELECT INTO} was executed
- \texttt{SPI\_OK\_INSERT}
  - if an \texttt{INSERT} was executed
SPI_OK_DELETE
   if a DELETE was executed
SPI_OK_UPDATE
   if an UPDATE was executed
SPI_OK_INSERT_RETURNING
   if an INSERT RETURNING was executed
SPI_OK_DELETE_RETURNING
   if a DELETE RETURNING was executed
SPI_OK_UPDATE_RETURNING
   if an UPDATE RETURNING was executed
SPI_OK Utility
   if a utility command (e.g., CREATE TABLE) was executed
SPI_OK_REWRITTEN
   if the command was rewritten into another kind of command (e.g., UPDATE became an INSERT) by a rule.

On error, one of the following negative values is returned:

SPI_ERROR_ARGUMENT
   if command is NULL or count is less than 0
SPI_ERROR_COPY
   if COPY TO stdout or COPY FROM stdin was attempted
SPI_ERROR_TRANSACTION
   if a transaction manipulation command was attempted (BEGIN, COMMIT, ROLLBACK, SAVE-POINT, PREPARE TRANSACTION, COMMIT PREPARED, ROLLBACK PREPARED, or any variant thereof)
SPI_ERROR_OPUNKNOWN
   if the command type is unknown (shouldn’t happen)
SPI_ERROR_UNCONNECTED
   if called from an unconnected C function

Notes

All SPI query-execution functions set both SPI_processed and SPI_tuptable (just the pointer, not the contents of the structure). Save these two global variables into local C function variables if you need to access the result table of SPI_execute or another query-execution function across later calls.
**SPI_exec**

SPI_exec — execute a read/write command

**Synopsis**

```c
int SPI_exec(const char * command, long count)
```

**Description**

SPI_exec is the same as SPI_execute, with the latter's read_only parameter always taken as false.

**Arguments**

- `const char * command`
  string containing command to execute
- `long count`
  maximum number of rows to return, or 0 for no limit

**Return Value**

See SPI_execute.
**SPI_execute_with_args**

SPI_execute_with_args — execute a command with out-of-line parameters

**Synopsis**

```c
int SPI_execute_with_args(const char *command,
                           int nargs, Oid *argtypes,
                           Datum *values, const char *nulls,
                           bool read_only, long count)
```

**Description**

SPI_execute_with_args executes a command that might include references to externally supplied parameters. The command text refers to a parameter as \$n, and the call specifies data types and values for each such symbol. read_only and count have the same interpretation as in SPI_execute.

The main advantage of this routine compared to SPI_execute is that data values can be inserted into the command without tedious quoting/escaping, and thus with much less risk of SQL-injection attacks.

Similar results can be achieved with SPI_prepare followed by SPI_execute_plan; however, when using this function the query plan is always customized to the specific parameter values provided. For one-time query execution, this function should be preferred. If the same command is to be executed with many different parameters, either method might be faster, depending on the cost of re-planning versus the benefit of custom plans.

**Arguments**

- const char * command
  - command string
- int nargs
  - number of input parameters ($1, $2, etc.)
- Oid * argtypes
  - an array of length nargs, containing the OIDs of the data types of the parameters
- Datum * values
  - an array of length nargs, containing the actual parameter values
- const char * nulls
  - an array of length nargs, describing which parameters are null

  If nulls is NULL then SPI_execute_with_args assumes that no parameters are null. Otherwise, each entry of the nulls array should be ' ' if the corresponding parameter value is non-null, or 'n' if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn't matter.) Note that nulls is not a text string, just an array: it does not need a '\0' terminator.
- bool read_only
  - true for read-only execution
long count

maximum number of rows to return, or 0 for no limit

Return Value

The return value is the same as for SPI_execute.

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.
**SPI_prepare**

SPI_prepare — prepare a statement, without executing it yet

**Synopsis**

```c
SPIPlanPtr SPI_prepare(const char * command, int nargs, Oid * argtypes)
```

**Description**

SPI_prepare creates and returns a prepared statement for the specified command, but doesn't execute the command. The prepared statement can later be executed repeatedly using SPI_execute_plan.

When the same or a similar command is to be executed repeatedly, it is generally advantageous to perform parse analysis only once, and might furthermore be advantageous to re-use an execution plan for the command. SPI_prepare converts a command string into a prepared statement that encapsulates the results of parse analysis. The prepared statement also provides a place for caching an execution plan if it is found that generating a custom plan for each execution is not helpful.

A prepared command can be generalized by writing parameters ($1, $2, etc.) in place of what would be constants in a normal command. The actual values of the parameters are then specified when SPI_execute_plan is called. This allows the prepared command to be used over a wider range of situations than would be possible without parameters.

The statement returned by SPI_prepare can be used only in the current invocation of the C function, since SPI_finish frees memory allocated for such a statement. But the statement can be saved for longer using the functions SPI_keepplan or SPI_saveplan.

**Arguments**

- `const char * command`
  - command string
- `int nargs`
  - number of input parameters ($1, $2, etc.)
- `Oid * argtypes`
  - pointer to an array containing the OIDs of the data types of the parameters

**Return Value**

SPI_prepare returns a non-null pointer to an SPIPlan, which is an opaque struct representing a prepared statement. On error, NULL will be returned, and SPI_result will be set to one of the same error codes used by SPI_execute, except that it is set to SPI_ERROR_ARGUMENT if command is NULL, or if nargs is less than 0, or if nargs is greater than 0 and argtypes is NULL.

**Notes**

If no parameters are defined, a generic plan will be created at the first use of SPI_execute_plan, and used for all subsequent executions as well. If there are parameters, the first few uses of SPI_execute_plan will generate custom plans that are specific to the supplied parameter values. Af-
ter enough uses of the same prepared statement, \texttt{SPI\_execute\_plan} will build a generic plan, and if that is not too much more expensive than the custom plans, it will start using the generic plan instead of re-planning each time. If this default behavior is unsuitable, you can alter it by passing the \texttt{CURSOR\_OPT\_GENERIC\_PLAN} or \texttt{CURSOR\_OPT\_CUSTOM\_PLAN} flag to \texttt{SPI\_prepare\_cursor}, to force use of generic or custom plans respectively.

Although the main point of a prepared statement is to avoid repeated parse analysis and planning of the statement, PostgreSQL will force re-analysis and re-planning of the statement before using it whenever database objects used in the statement have undergone definitional (DDL) changes since the previous use of the prepared statement. Also, if the value of search\_path changes from one use to the next, the statement will be re-parsed using the new search\_path. (This latter behavior is new as of PostgreSQL 9.3.) See \texttt{PREPARE} for more information about the behavior of prepared statements.

This function should only be called from a connected C function.

\texttt{SPIPlanPtr} is declared as a pointer to an opaque struct type in \texttt{spi.h}. It is unwise to try to access its contents directly, as that makes your code much more likely to break in future revisions of PostgreSQL.

The name \texttt{SPIPlanPtr} is somewhat historical, since the data structure no longer necessarily contains an execution plan.
SPI_prepare_cursor

SPI_prepare_cursor — prepare a statement, without executing it yet

Synopsis

SPIPlanPtr SPI_prepare_cursor(const char * command, int nargs,
                                Oid * argtypes, int cursorOptions)

Description

SPI_prepare_cursor is identical to SPI_prepare, except that it also allows specification of
the planner's “cursor options” parameter. This is a bit mask having the values shown in nodes/
parsenodes.h for the options field of DeclareCursorStmt. SPI_prepare always takes
the cursor options as zero.

Arguments

const char * command
    command string
int nargs
    number of input parameters ($1, $2, etc.)
Oid * argtypes
    pointer to an array containing the OIDs of the data types of the parameters
int cursorOptions
    integer bit mask of cursor options; zero produces default behavior

Return Value

SPI_prepare_cursor has the same return conventions as SPI_prepare.

Notes

Useful bits to set in cursorOptions include CURSOR_OPT_SCROLL,
CURSOR_OPT_NO_SCROLL, CURSOR_OPT_FAST_PLAN, CURSOR_OPT_GENERIC_PLAN,
and CURSOR_OPT_CUSTOM_PLAN. Note in particular that CURSOR_OPT_HOLD is ignored.
**SPI_prepare_params**

SPI_prepare_params — prepare a statement, without executing it yet

**Synopsis**

```c
SPIPlanPtr SPI_prepare_params(const char * command,
                                ParserSetupHook parserSetup,
                                void * parserSetupArg,
                                int cursorOptions)
```

**Description**

SPI_prepare_params creates and returns a prepared statement for the specified command, but doesn’t execute the command. This function is equivalent to SPI_prepare_cursor, with the addition that the caller can specify parser hook functions to control the parsing of external parameter references.

**Arguments**

- `const char * command`
  - command string
- `ParserSetupHook parserSetup`
  - Parser hook setup function
- `void * parserSetupArg`
  - pass-through argument for `parserSetup`
- `int cursorOptions`
  - integer bit mask of cursor options; zero produces default behavior

**Return Value**

SPI_prepare_params has the same return conventions as SPI_prepare.
**SPI_getargcount**

SPI_getargcount — return the number of arguments needed by a statement prepared by SPI_prepare

**Synopsis**

```c
int SPI_getargcount(SPIPlanPtr plan)
```

**Description**

SPI_getargcount returns the number of arguments needed to execute a statement prepared by SPI_prepare.

**Arguments**

- `SPIPlanPtr plan`
  
  prepared statement (returned by SPI_prepare)

**Return Value**

The count of expected arguments for the `plan`. If the `plan` is NULL or invalid, SPI_result is set to SPI_ERROR_ARGUMENT and -1 is returned.
SPI_getargtypeid

SPI_getargtypeid — return the data type OID for an argument of a statement prepared by SPI_prepare

Synopsis

Oid SPI_getargtypeid(SPIPlanPtr plan, int argIndex)

Description

SPI_getargtypeid returns the OID representing the type for the argIndex' th argument of a statement prepared by SPI_prepare. First argument is at index zero.

Arguments

SPIPlanPtr plan

prepared statement (returned by SPI_prepare)

int argIndex

zero based index of the argument

Return Value

The type OID of the argument at the given index. If the plan is NULL or invalid, or argIndex is less than 0 or not less than the number of arguments declared for the plan, SPI_result is set to SPI_ERROR_ARGUMENT and InvalidOid is returned.
SPI_is_cursor_plan

SPI_is_cursor_plan — return true if a statement prepared by SPI_prepare can be used with SPI_cursor_open

Synopsis

bool SPI_is_cursor_plan(SPIPlanPtr plan)

Description

SPI_is_cursor_plan returns true if a statement prepared by SPI_prepare can be passed as an argument to SPI_cursor_open, or false if that is not the case. The criteria are that the plan represents one single command and that this command returns tuples to the caller; for example, SELECT is allowed unless it contains an INTO clause, and UPDATE is allowed only if it contains a RETURNING clause.

Arguments

SPIPlanPtr plan

prepared statement (returned by SPI_prepare)

Return Value

ture or false to indicate if the plan can produce a cursor or not, with SPI_result set to zero. If it is not possible to determine the answer (for example, if the plan is NULL or invalid, or if called when not connected to SPI), then SPI_result is set to a suitable error code and false is returned.
SPI_execute_plan

SPI_execute_plan — execute a statement prepared by SPI_prepare

Synopsis

    int SPI_execute_plan(SPIPlanPtr plan, Datum * values, const char * nulls,
                        bool read_only, long count)

Description

SPI_execute_plan executes a statement prepared by SPI_prepare or one of its siblings. read_only and count have the same interpretation as in SPI_execute.

Arguments

SPIPlanPtr plan

    prepared statement (returned by SPI_prepare)

Datum * values

    An array of actual parameter values. Must have same length as the statement's number of arguments.

const char * nulls

    An array describing which parameters are null. Must have same length as the statement's number of arguments.

    If nulls is NULL then SPI_execute_plan assumes that no parameters are null. Otherwise, each entry of the nulls array should be ' ' if the corresponding parameter value is non-null, or 'n' if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn't matter.) Note that nulls is not a text string, just an array: it does not need a '\0' terminator.

bool read_only

    true for read-only execution

long count

    maximum number of rows to return, or 0 for no limit

Return Value

The return value is the same as for SPI_execute, with the following additional possible error (negative) results:

SPI_ERROR_ARGUMENT

    if plan is NULL or invalid, or count is less than 0

SPI_ERROR_PARAM

    if values is NULL and plan was prepared with some parameters
SPI_processed and SPI_tuptable are set as in SPI_execute if successful.
**SPI_execute_plan_with_paramlist**

SPI_execute_plan_with_paramlist — execute a statement prepared by SPI_prepare

**Synopsis**

```c
int SPI_execute_plan_with_paramlist(SPIPlanPtr plan,
                                    ParamListInfo params,
                                    bool read_only,
                                    long count)
```

**Description**

SPI_execute_plan_with_paramlist executes a statement prepared by SPI_prepare. This function is equivalent to SPI_execute_plan except that information about the parameter values to be passed to the query is presented differently. The ParamListInfo representation can be convenient for passing down values that are already available in that format. It also supports use of dynamic parameter sets via hook functions specified in ParamListInfo.

**Arguments**

- **SPIPlanPtr plan**
  - prepared statement (returned by SPI_prepare)
- **ParamListInfo params**
  - data structure containing parameter types and values; NULL if none
- **bool read_only**
  - true for read-only execution
- **long count**
  - maximum number of rows to return, or 0 for no limit

**Return Value**

The return value is the same as for SPI_execute_plan.

SPI_processed and SPI_tuptable are set as in SPI_execute_plan if successful.
SPI_execp

SPI_execp — execute a statement in read/write mode

Synopsis

```
int SPI_execp(SPIPlanPtr plan, Datum * values, const char * nulls,
             long count)
```

Description

SPI_execp is the same as SPI_execute_plan, with the latter's read_only parameter always taken as false.

Arguments

SPIPlanPtr plan

preparedStatement (returned by SPI_prepare)

Datum * values

An array of actual parameter values. Must have same length as the statement's number of arguments.

const char * nulls

An array describing which parameters are null. Must have same length as the statement's number of arguments.

If nulls is NULL then SPI_execp assumes that no parameters are null. Otherwise, each entry of the nulls array should be ' ' if the corresponding parameter value is non-null, or 'n' if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn't matter.) Note that nulls is not a text string, just an array: it does not need a '\0' terminator.

long count

maximum number of rows to return, or 0 for no limit

Return Value

See SPI_execute_plan.

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.
**SPI_cursor_open**

SPI_cursor_open — set up a cursor using a statement created with SPI_prepare

**Synopsis**

```c
Portal SPI_cursor_open(const char * name, SPIPlanPtr plan,
    Datum * values, const char * nulls,
    bool read_only)
```

**Description**

SPI_cursor_open sets up a cursor (internally, a portal) that will execute a statement prepared by SPI_prepare. The parameters have the same meanings as the corresponding parameters to SPI_execute_plan.

Using a cursor instead of executing the statement directly has two benefits. First, the result rows can be retrieved a few at a time, avoiding memory overrun for queries that return many rows. Second, a portal can outlive the current C function (it can, in fact, live to the end of the current transaction). Returning the portal name to the C function's caller provides a way of returning a row set as result.

The passed-in parameter data will be copied into the cursor's portal, so it can be freed while the cursor still exists.

**Arguments**

- `const char * name`
  
  name for portal, or NULL to let the system select a name

- `SPIPlanPtr plan`
  
  prepared statement (returned by SPI_prepare)

- `Datum * values`
  
  An array of actual parameter values. Must have same length as the statement's number of arguments.

- `const char * nulls`
  
  An array describing which parameters are null. Must have same length as the statement's number of arguments.

  If `nulls` is NULL then SPI_cursor_open assumes that no parameters are null. Otherwise, each entry of the `nulls` array should be ' ' if the corresponding parameter value is non-null, or 'n' if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding `values` entry doesn't matter.) Note that `nulls` is not a text string, just an array: it does not need a '\0' terminator.

- `bool read_only`
  
  true for read-only execution

**Return Value**

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via `elog`. 
SPI_cursor_open_with_args

SPI_cursor_open_with_args — set up a cursor using a query and parameters

Synopsis

Portal SPI_cursor_open_with_args(const char *name,
        const char *command,
        int nargs, Oid *argtypes,
        Datum *values, const char *nulls,
        bool read_only, int cursorOptions)

Description

SPI_cursor_open_with_args sets up a cursor (internally, a portal) that will execute the specified query. Most of the parameters have the same meanings as the corresponding parameters to SPI_prepare_cursor and SPI_cursor_open.

For one-time query execution, this function should be preferred over SPI_prepare_cursor followed by SPI_cursor_open. If the same command is to be executed with many different parameters, either method might be faster, depending on the cost of re-planning versus the benefit of custom plans.

The passed-in parameter data will be copied into the cursor's portal, so it can be freed while the cursor still exists.

Arguments

cost char * name
   name for portal, or NULL to let the system select a name

cost char * command
   command string

int nargs
   number of input parameters ($1, $2, etc.)

Oid * argtypes
   an array of length nargs, containing the OIDs of the data types of the parameters

Datum * values
   an array of length nargs, containing the actual parameter values

const char * nulls
   an array of length nargs, describing which parameters are null

If nulls is NULL then SPI_cursor_open_with_args assumes that no parameters are null. Otherwise, each entry of the nulls array should be '' if the corresponding parameter value is non-null, or 'n' if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn't matter.) Note that nulls is not a text string, just an array: it does not need a '\0' terminator.
bool read_only

true for read-only execution

int cursorOptions

integer bit mask of cursor options; zero produces default behavior

**Return Value**

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via `elog`. 
**SPI_cursor_open_with_paramlist**

SPI_cursor_open_with_paramlist — set up a cursor using parameters

**Synopsis**

```c
Portal SPI_cursor_open_with_paramlist(const char *name,
SPIPlanPtr plan,
ParamListInfo params,
bool read_only)
```

**Description**

SPI_cursor_open_with_paramlist sets up a cursor (internally, a portal) that will execute a statement prepared by SPI_prepare. This function is equivalent to SPI_cursor_open except that information about the parameter values to be passed to the query is presented differently. The ParamListInfo representation can be convenient for passing down values that are already available in that format. It also supports use of dynamic parameter sets via hook functions specified in ParamListInfo.

The passed-in parameter data will be copied into the cursor's portal, so it can be freed while the cursor still exists.

**Arguments**

- `const char * name`
  - name for portal, or NULL to let the system select a name
- `SPIPlanPtr plan`
  - prepared statement (returned by SPI_prepare)
- `ParamListInfo params`
  - data structure containing parameter types and values; NULL if none
- `bool read_only`
  - true for read-only execution

**Return Value**

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via elog.
SPI_cursor_find

SPI_cursor_find — find an existing cursor by name

Synopsis

Portal SPI_cursor_find(const char * name)

Description

SPI_cursor_find finds an existing portal by name. This is primarily useful to resolve a cursor name returned as text by some other function.

Arguments

const char * name

name of the portal

Return Value

pointer to the portal with the specified name, or NULL if none was found
SPI_cursor_fetch

SPI_cursor_fetch — fetch some rows from a cursor

Synopsis

void SPI_cursor_fetch(Portal portal, bool forward, long count)

Description

SPI_cursor_fetch fetches some rows from a cursor. This is equivalent to a subset of the SQL command FETCH (see SPI_scroll_cursor_fetch for more functionality).

Arguments

Portal portal
  portal containing the cursor
bool forward
  true for fetch forward, false for fetch backward
long count
  maximum number of rows to fetch

Return Value

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.

Notes

Fetching backward may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
**SPI_cursor_move**

SPI_cursor_move — move a cursor

**Synopsis**

```c
void SPI_cursor_move(Portal portal, bool forward, long count)
```

**Description**

SPI_cursor_move skips over some number of rows in a cursor. This is equivalent to a subset of the SQL command MOVE (see SPI_scroll_cursor_move for more functionality).

**Arguments**

- **Portal portal**
  - portal containing the cursor
- **bool forward**
  - true for move forward, false for move backward
- **long count**
  - maximum number of rows to move

**Notes**

Moving backward may fail if the cursor's plan was not created with the CURSOR_OPT_SCROLL option.
**SPI_scroll_cursor_fetch**

SPI_scroll_cursor_fetch — fetch some rows from a cursor

**Synopsis**

```c
void SPI_scroll_cursor_fetch(Portal portal, FetchDirection direction, long count)
```

**Description**

SPI_scroll_cursor_fetch fetches some rows from a cursor. This is equivalent to the SQL command FETCH.

**Arguments**

- `Portal portal` - portal containing the cursor
- `FetchDirection direction` - one of FETCH_FORWARD, FETCH_BACKWARD, FETCH_ABSOLUTE or FETCH_RELATIVE
- `long count` - number of rows to fetch for FETCH_FORWARD or FETCH_BACKWARD; absolute row number to fetch for FETCH_ABSOLUTE; or relative row number to fetch for FETCH_RELATIVE

**Return Value**

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.

**Notes**

See the SQL FETCH command for details of the interpretation of the `direction` and `count` parameters.

Direction values other than FETCH_FORWARD may fail if the cursor's plan was not created with the CURSOR_OPT_SCROLL option.
**SPI_scroll_cursor_move**

SPI_scroll_cursor_move — move a cursor

**Synopsis**

```c
void SPI_scroll_cursor_move(Portal portal,
    FetchDirection direction,
    long count)
```

**Description**

SPI_scroll_cursor_move skips over some number of rows in a cursor. This is equivalent to the SQL command MOVE.

**Arguments**

- **Portal portal**
  
  portal containing the cursor
- **FetchDirection direction**
  
  one of FETCH_FORWARD, FETCH_BACKWARD, FETCH_ABSOLUTE or FETCH_RELATIVE
- **long count**
  
  number of rows to move for FETCH_FORWARD or FETCH_BACKWARD; absolute row number to move to for FETCH_ABSOLUTE; or relative row number to move to for FETCH_RELATIVE

**Return Value**

SPI_processed is set as in SPI_execute if successful. SPI_tuptable is set to NULL, since no rows are returned by this function.

**Notes**

See the SQL FETCH command for details of the interpretation of the direction and count parameters.

Direction values other than FETCH_FORWARD may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
SPI_cursor_close

SPI_cursor_close — close a cursor

Synopsis

```c
void SPI_cursor_close(Portal portal)
```

Description

SPI_cursor_close closes a previously created cursor and releases its portal storage.

All open cursors are closed automatically at the end of a transaction. SPI_cursor_close need only be invoked if it is desirable to release resources sooner.

Arguments

- **Portal portal**
  
  portal containing the cursor
SPI_keepplan

SPI_keepplan — save a prepared statement

Synopsis

```c
int SPI_keepplan(SPIPlanPtr plan)
```

Description

SPI_keepplan saves a passed statement (prepared by SPI_prepare) so that it will not be freed by SPI_finish nor by the transaction manager. This gives you the ability to reuse prepared statements in the subsequent invocations of your C function in the current session.

Arguments

```c
SPIPlanPtr plan
```

the prepared statement to be saved

Return Value

0 on success; SPI_ERROR_ARGUMENT if `plan` is NULL or invalid

Notes

The passed-in statement is relocated to permanent storage by means of pointer adjustment (no data copying is required). If you later wish to delete it, use SPI_freeplan on it.
**SPI_saveplan**

SPI_saveplan — save a prepared statement

**Synopsis**

```c
SPIPlanPtr SPI_saveplan(SPIPlanPtr plan)
```

**Description**

SPI_saveplan copies a passed statement (prepared by SPI_prepare) into memory that will not be freed by SPI_finish nor by the transaction manager, and returns a pointer to the copied statement. This gives you the ability to reuse prepared statements in the subsequent invocations of your C function in the current session.

**Arguments**

- `SPIPlanPtr plan`
  - the prepared statement to be saved

**Return Value**

- Pointer to the copied statement; or NULL if unsuccessful. On error, SPI_result is set thus:
  - **SPI_ERROR_ARGUMENT**
    - if `plan` is NULL or invalid
  - **SPI_ERROR_UNCONNECTED**
    - if called from an unconnected C function

**Notes**

The originally passed-in statement is not freed, so you might wish to do SPI_freeplan on it to avoid leaking memory until SPI_finish.

In most cases, SPI_keepplan is preferred to this function, since it accomplishes largely the same result without needing to physically copy the prepared statement's data structures.
SPI_register_relation

SPI_register_relation — make an ephemeral named relation available by name in SPI queries

Synopsis

int SPI_register_relation(EphemeralNamedRelation enr)

Description

SPI_register_relation makes an ephemeral named relation, with associated information, available to queries planned and executed through the current SPI connection.

Arguments

EphemeralNamedRelation enr

the ephemeral named relation registry entry

Return Value

If the execution of the command was successful then the following (nonnegative) value will be returned:

SPI_OK_REL_REGISTER

if the relation has been successfully registered by name

On error, one of the following negative values is returned:

SPI_ERROR_ARGUMENT

if enr is NULL or its name field is NULL

SPI_ERROR_UNCONNECTED

if called from an unconnected C function

SPI_ERROR_REL_DUPLICATE

if the name specified in the name field of enr is already registered for this connection
SPI_unregister_relation

SPI_unregister_relation — remove an ephemeral named relation from the registry

Synopsis

int SPI_unregister_relation(const char * name)

Description

SPI_unregister_relation removes an ephemeral named relation from the registry for the current connection.

Arguments

const char * name

the relation registry entry name

Return Value

If the execution of the command was successful then the following (nonnegative) value will be returned:

SPI_OK_REL_UNREGISTER

if the tuplestore has been successfully removed from the registry

On error, one of the following negative values is returned:

SPI_ERROR_ARGUMENT

if name is NULL

SPI_ERROR_UNCONNECTED

if called from an unconnected C function

SPI_ERROR_REL_NOT_FOUND

if name is not found in the registry for the current connection
**SPI_register_trigger_data**

SPI_register_trigger_data — make ephemeral trigger data available in SPI queries

### Synopsis

```c
int SPI_register_trigger_data(TriggerData *tdata)
```

### Description

SPI_register_trigger_data makes any ephemeral relations captured by a trigger available to queries planned and executed through the current SPI connection. Currently, this means the transition tables captured by an *AFTER* trigger defined with a *REFERENCING OLD/New TABLE AS* clause. This function should be called by a PL trigger handler function after connecting.

### Arguments

- **TriggerData *tdata**
  
  the TriggerData object passed to a trigger handler function as `fcinfo->context`

### Return Value

If the execution of the command was successful then the following (nonnegative) value will be returned:

- **SPI_OK_TD_REGISTER**
  
  if the captured trigger data (if any) has been successfully registered

On error, one of the following negative values is returned:

- **SPI_ERROR_ARGUMENT**
  
  if `tdata` is NULL

- **SPI_ERROR_UNCONNECTED**
  
  if called from an unconnected C function

- **SPI_ERROR_REL_DUPLICATE**
  
  if the name of any trigger data transient relation is already registered for this connection

### 47.2. Interface Support Functions

The functions described here provide an interface for extracting information from result sets returned by `SPI_execute` and other SPI functions.

All functions described in this section can be used by both connected and unconnected C functions.
SPI_fname

SPI_fname — determine the column name for the specified column number

Synopsis

char * SPI_fname(TupleDesc rowdesc, int colnumber)

Description

SPI_fname returns a copy of the column name of the specified column. (You can use pfree to release the copy of the name when you don't need it anymore.)

Arguments

TupleDesc rowdesc
   input row description
int colnumber
   column number (count starts at 1)

Return Value

The column name; NULL if colnumber is out of range. SPI_result set to SPI_ERROR_NOATTRIBUTE on error.
SPI_fnumber

SPI_fnumber — determine the column number for the specified column name

Synopsis

```c
int SPI_fnumber(TupleDesc rowdesc, const char * colname)
```

Description

SPI_fnumber returns the column number for the column with the specified name.

If `colname` refers to a system column (e.g., `oid`) then the appropriate negative column number will be returned. The caller should be careful to test the return value for exact equality to `SPI_ERROR_NOATTRIBUTE` to detect an error; testing the result for less than or equal to 0 is not correct unless system columns should be rejected.

Arguments

- `TupleDesc rowdesc`  
  input row description
- `const char * colname`  
  column name

Return Value

Column number (count starts at 1 for user-defined columns), or `SPI_ERROR_NOATTRIBUTE` if the named column was not found.
SPI_getvalue

SPI_getvalue — return the string value of the specified column

Synopsis

char * SPI_getvalue(HeapTuple row, TupleDesc rowdesc,
int colnumber)

Description

SPI_getvalue returns the string representation of the value of the specified column.

The result is returned in memory allocated using palloc. (You can use pfree to release the memory when you don’t need it anymore.)

Arguments

HeapTuple row

input row to be examined

TupleDesc rowdesc

input row description

int colnumber

column number (count starts at 1)

Return Value

Column value, or NULL if the column is null, colnumber is out of range (SPI_result is set to SPI_ERROR_NOATTRIBUTE), or no output function is available (SPI_result is set to SPI_ERROR_NOOUTFUNC).
SPI_getbinval

SPI_getbinval — return the binary value of the specified column

Synopsis

Datum SPI_getbinval(HeapTuple row, TupleDesc rowdesc, int colnumber, bool * isnull)

Description

SPI_getbinval returns the value of the specified column in the internal form (as type Datum).

This function does not allocate new space for the datum. In the case of a pass-by-reference data type, the return value will be a pointer into the passed row.

Arguments

HeapTuple row
input row to be examined

TupleDesc rowdesc
input row description

int colnumber
column number (count starts at 1)

bool * isnull
flag for a null value in the column

Return Value

The binary value of the column is returned. The variable pointed to by isnull is set to true if the column is null, else to false.

SPI_result is set to SPI_ERROR_NOATTRIBUTE on error.
SPI_gettype

SPI_gettype — return the data type name of the specified column

Synopsis

char * SPI_gettype(TupleDesc rowdesc, int colnumber)

Description

SPI_gettype returns a copy of the data type name of the specified column. (You can use pfree to release the copy of the name when you don’t need it anymore.)

Arguments

TupleDesc rowdesc
input row description

int colnumber
column number (count starts at 1)

Return Value

The data type name of the specified column, or NULL on error. SPI_result is set to SPI_ERROR_NOATTRIBUTE on error.
SPI_gettypeid

SPI_gettypeid — return the data type OID of the specified column

Synopsis

Oid SPI_gettypeid(TupleDesc rowdesc, int colnumber)

Description

SPI_gettypeid returns the OID of the data type of the specified column.

Arguments

TupleDesc rowdesc

input row description

int colnumber

column number (count starts at 1)

Return Value

The OID of the data type of the specified column or InvalidOid on error. On error, SPI_result is set to SPI_ERROR_NOATTRIBUTE.
SPI_getrelname

SPI_getrelname — return the name of the specified relation

Synopsis

char * SPI_getrelname(Relation rel)

Description

SPI_getrelname returns a copy of the name of the specified relation. (You can use pfree to release the copy of the name when you don't need it anymore.)

Arguments

Relation rel

input relation

Return Value

The name of the specified relation.
SPI_getnspname

SPI_getnspname — return the namespace of the specified relation

Synopsis

char * SPI_getnspname(Relation rel)

Description

SPI_getnspname returns a copy of the name of the namespace that the specified Relation belongs to. This is equivalent to the relation's schema. You should pfree the return value of this function when you are finished with it.

Arguments

Relation rel

input relation

Return Value

The name of the specified relation's namespace.
**SPI_result_code_string**

SPI_result_code_string — return error code as string

**Synopsis**

```c
const char * SPI_result_code_string(int code);
```

**Description**

`SPI_result_code_string` returns a string representation of the result code returned by various SPI functions or stored in `SPI_result`.

**Arguments**

- `int code`
  - result code

**Return Value**

A string representation of the result code.

### 47.3. Memory Management

PostgreSQL allocates memory within *memory contexts*, which provide a convenient method of managing allocations made in many different places that need to live for differing amounts of time. Destroying a context releases all the memory that was allocated in it. Thus, it is not necessary to keep track of individual objects to avoid memory leaks; instead only a relatively small number of contexts have to be managed. `palloc` and related functions allocate memory from the “current” context.

`SPI_connect` creates a new memory context and makes it current. `SPI_finish` restores the previous current memory context and destroys the context created by `SPI_connect`. These actions ensure that transient memory allocations made inside your C function are reclaimed at C function exit, avoiding memory leakage.

However, if your C function needs to return an object in allocated memory (such as a value of a pass-by-reference data type), you cannot allocate that memory using `palloc`, at least not while you are connected to SPI. If you try, the object will be deallocated by `SPI_finish`, and your C function will not work reliably. To solve this problem, use `SPI_palloc` to allocate memory for your return object. `SPI_palloc` allocates memory in the “upper executor context”, that is, the memory context that was current when `SPI_connect` was called, which is precisely the right context for a value returned from your C function. Several of the other utility functions described in this section also return objects created in the upper executor context.

When `SPI_connect` is called, the private context of the C function, which is created by `SPI_connect`, is made the current context. All allocations made by `palloc`, `repalloc`, or SPI utility functions (except as described in this section) are made in this context. When a C function disconnects from the SPI manager (via `SPI_finish`) the current context is restored to the upper executor context, and all allocations made in the C function memory context are freed and cannot be used any more.
SPI_palloc

SPI_palloc — allocate memory in the upper executor context

Synopsis

void * SPI_palloc(Size size)

Description

SPI_palloc allocates memory in the upper executor context.

This function can only be used while connected to SPI. Otherwise, it throws an error.

Arguments

Size size

size in bytes of storage to allocate

Return Value

pointer to new storage space of the specified size
SPI_realloc

SPI_realloc — reallocate memory in the upper executor context

Synopsis

void * SPI_realloc(void * pointer, Size size)

Description

SPI_realloc changes the size of a memory segment previously allocated using SPI_palloc. This function is no longer different from plain repalloc. It's kept just for backward compatibility of existing code.

Arguments

void * pointer

pointer to existing storage to change

Size size

size in bytes of storage to allocate

Return Value

pointer to new storage space of specified size with the contents copied from the existing area
**SPI_pfree**

SPI_pfree — free memory in the upper executor context

**Synopsis**

```c
void SPI_pfree(void * pointer)
```

**Description**

SPI_pfree frees memory previously allocated using SPI_palloc or SPI_repalloc.

This function is no longer different from plain `pfree`. It's kept just for backward compatibility of existing code.

**Arguments**

```c
void * pointer
```

pointer to existing storage to free
**SPI_copytuple**

SPI_copytuple — make a copy of a row in the upper executor context

**Synopsis**

HeapTuple SPI_copytuple(HeapTuple row)

**Description**

SPI_copytuple makes a copy of a row in the upper executor context. This is normally used to return a modified row from a trigger. In a function declared to return a composite type, use SPI_returntuple instead.

This function can only be used while connected to SPI. Otherwise, it returns NULL and sets SPI_result to SPI_ERROR_UNCONNECTED.

**Arguments**

HeapTuple row

row to be copied

**Return Value**

the copied row, or NULL on error (see SPI_result for an error indication)
**SPI_returntuple**

SPI_returntuple — prepare to return a tuple as a Datum

**Synopsis**

```
HeapTupleHeader SPI_returntuple(HeapTuple row, TupleDesc rowdesc)
```

**Description**

SPI_returntuple makes a copy of a row in the upper executor context, returning it in the form of a row type Datum. The returned pointer need only be converted to Datum via PointerGetDatum before returning.

This function can only be used while connected to SPI. Otherwise, it returns NULL and sets SPI_result to SPI_ERROR_UNCONNECTED.

Note that this should be used for functions that are declared to return composite types. It is not used for triggers; use SPI_copytuple for returning a modified row in a trigger.

**Arguments**

- **HeapTuple row**
  - row to be copied
- **TupleDesc rowdesc**
  - descriptor for row (pass the same descriptor each time for most effective caching)

**Return Value**

- `HeapTupleHeader` pointing to copied row, or NULL on error (see SPI_result for an error indication)
**SPI_modifytuple**

SPI_modifytuple — create a row by replacing selected fields of a given row

**Synopsis**

```c
HeapTuple SPI_modifytuple(Relation rel, HeapTuple row, int ncols, int * colnum, Datum * values, const char * nulls)
```

**Description**

SPI_modifytuple creates a new row by substituting new values for selected columns, copying the original row's columns at other positions. The input row is not modified. The new row is returned in the upper executor context.

This function can only be used while connected to SPI. Otherwise, it returns NULL and sets SPI_result to SPI_ERROR_UNCONNECTED.

**Arguments**

- **Relation rel**
  
  Used only as the source of the row descriptor for the row. (Passing a relation rather than a row descriptor is a misfeature.)

- **HeapTuple row**
  
  row to be modified

- **int ncols**
  
  number of columns to be changed

- **int * colnum**
  
  an array of length ncols, containing the numbers of the columns that are to be changed (column numbers start at 1)

- **Datum * values**
  
  an array of length ncols, containing the new values for the specified columns

- **const char * nulls**
  
  an array of length ncols, describing which new values are null

  If nulls is NULL then SPI_modifytuple assumes that no new values are null. Otherwise, each entry of the nulls array should be ' ' if the corresponding new value is non-null, or 'n' if the corresponding new value is null. (In the latter case, the actual value in the corresponding values entry doesn't matter.) Note that nulls is not a text string, just an array: it does not need a '\0' terminator.

**Return Value**

new row with modifications, allocated in the upper executor context, or NULL on error (see SPI_result for an error indication)
On error, SPI_result is set as follows:

SPI_ERROR_ARGUMENT

if rel is NULL, or if row is NULL, or if ncols is less than or equal to 0, or if colnum is NULL, or if values is NULL.

SPI_ERROR_NOATTRIBUTE

if colnum contains an invalid column number (less than or equal to 0 or greater than the number of columns in row)

SPI_ERROR_UNCONNECTED

if SPI is not active
**SPI_freetuple**

SPI_freetuple — free a row allocated in the upper executor context

**Synopsis**

```c
void SPI_freetuple(HeapTuple row)
```

**Description**

SPI_freetuple frees a row previously allocated in the upper executor context.

This function is no longer different from plain heap_freetuple. It's kept just for backward compatibility of existing code.

**Arguments**

- **HeapTuple row**
  - row to free
**SPI_freetuptable**

SPI_freetuptable — free a row set created by SPI_execute or a similar function

**Synopsis**

```c
void SPI_freetuptable(SPITupleTable * tuptable)
```

**Description**

SPI_freetuptable frees a row set created by a prior SPI command execution function, such as SPI_execute. Therefore, this function is often called with the global variable SPI_tuptable as argument.

This function is useful if an SPI-using C function needs to execute multiple commands and does not want to keep the results of earlier commands around until it ends. Note that any unfreed row sets will be freed anyway at SPI_finish. Also, if a subtransaction is started and then aborted within execution of an SPI-using C function, SPI automatically frees any row sets created while the subtransaction was running.

Beginning in PostgreSQL 9.3, SPI_freetuptable contains guard logic to protect against duplicate deletion requests for the same row set. In previous releases, duplicate deletions would lead to crashes.

**Arguments**

```c
SPITupleTable * tuptable
```

pointer to row set to free, or NULL to do nothing
**SPI_freeplan**

SPI_freeplan — free a previously saved prepared statement

**Synopsis**

```c
int SPI_freeplan(SPIPlanPtr plan)
```

**Description**

SPI_freeplan releases a prepared statement previously returned by SPI_prepare or saved by SPI_keepplan or SPI_saveplan.

**Arguments**

- `SPIPlanPtr plan`

  pointer to statement to free

**Return Value**

- 0 on success; SPI_ERROR_ARGUMENT if `plan` is NULL or invalid

---

**47.4. Transaction Management**

It is not possible to run transaction control commands such as COMMIT and ROLLBACK through SPI functions such as SPI_execute. There are, however, separate interface functions that allow transaction control through SPI.

It is not generally safe and sensible to start and end transactions in arbitrary user-defined SQL-callable functions without taking into account the context in which they are called. For example, a transaction boundary in the middle of a function that is part of a complex SQL expression that is part of some SQL command will probably result in obscure internal errors or crashes. The interface functions presented here are primarily intended to be used by procedural language implementations to support transaction management in SQL-level procedures that are invoked by the CALL command, taking the context of the CALL invocation into account. SPI-using procedures implemented in C can implement the same logic, but the details of that are beyond the scope of this documentation.
SPI_commit

SPI_commit — commit the current transaction

Synopsis

void SPI_commit(void)

Description

SPI_commit commits the current transaction. It is approximately equivalent to running the SQL command COMMIT. After a transaction is committed, a new transaction has to be started using SPI_start_transaction before further database actions can be executed.

This function can only be executed if the SPI connection has been set as nonatomic in the call to SPI_connect_ext.
**SPI_rollback**

SPI_rollback — abort the current transaction

**Synopsis**

```c
void SPI_rollback(void)
```

**Description**

SPI_rollback rolls back the current transaction. It is approximately equivalent to running the SQL command ROLLBACK. After a transaction is rolled back, a new transaction has to be started using SPI_start_transaction before further database actions can be executed.

This function can only be executed if the SPI connection has been set as nonatomic in the call to SPI_connect_ext.
**SPI_start_transaction**

SPI_start_transaction — start a new transaction

**Synopsis**

```c
void SPI_start_transaction(void)
```

**Description**

SPI_start_transaction starts a new transaction. It can only be called after SPI_commit or SPI_rollback, as there is no transaction active at that point. Normally, when an SPI-using procedure is called, there is already a transaction active, so attempting to start another one before closing out the current one will result in an error.

This function can only be executed if the SPI connection has been set as nonatomic in the call to SPI_connect_ext.

### 47.5. Visibility of Data Changes

The following rules govern the visibility of data changes in functions that use SPI (or any other C function):

- During the execution of an SQL command, any data changes made by the command are invisible to the command itself. For example, in:

  ```sql
  INSERT INTO a SELECT * FROM a;
  ```

  the inserted rows are invisible to the SELECT part.

- Changes made by a command C are visible to all commands that are started after C, no matter whether they are started inside C (during the execution of C) or after C is done.

- Commands executed via SPI inside a function called by an SQL command (either an ordinary function or a trigger) follow one or the other of the above rules depending on the read/write flag passed to SPI. Commands executed in read-only mode follow the first rule: they cannot see changes of the calling command. Commands executed in read-write mode follow the second rule: they can see all changes made so far.

- All standard procedural languages set the SPI read-write mode depending on the volatility attribute of the function. Commands of STABLE and IMMUTABLE functions are done in read-only mode, while commands of VOLATILE functions are done in read-write mode. While authors of C functions are able to violate this convention, it's unlikely to be a good idea to do so.

The next section contains an example that illustrates the application of these rules.

### 47.6. Examples

This section contains a very simple example of SPI usage. The C function `execq` takes an SQL command as its first argument and a row count as its second, executes the command using SPI_exec and returns the number of rows that were processed by the command. You can find more complex examples for SPI in the source tree in `src/test/regress/regress.c` and in the `spi` module.

```c
#include "postgres.h"
```
Server Programming Interface

#include "executor/spi.h"
#include "utils/builtins.h"

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(execq);

Datum
execq(PG_FUNCTION_ARGS)
{
  char *command;
  int cnt;
  int ret;
  uint64 proc;

  /* Convert given text object to a C string */
  command = text_to_cstring(PG_GETARG_TEXT_PP(1));
  cnt = PG_GETARG_INT32(2);

  SPI_connect();

  ret = SPI_exec(command, cnt);

  proc = SPI_processed;

  /* If some rows were fetched, print them via elog(INFO). */
  if (ret > 0 && SPI_tuptable != NULL)
  {
    TupleDesc tupdesc = SPI_tuptable->tupdesc;
    SPITupleTable *tuptable = SPI_tuptable;
    char buf[8192];
    uint64 j;

    for (j = 0; j < proc; j++)
    {
      HeapTuple tuple = tuptable->vals[j];
      int i;

      for (i = 1, buf[0] = 0; i <= tupdesc->natts; i++)
        snprintf(buf + strlen(buf), sizeof(buf) -
                  strlen(buf), "%s%s",
                  SPI_getvalue(tuple, tupdesc, i),
                  (i == tupdesc->natts) ? " " : " |");
      elog(INFO, "EXECQ: %s", buf);
    }
  }

  SPI_finish();
  pfree(command);

  PG_RETURN_INT64(proc);
}

This is how you declare the function after having compiled it into a shared library (details are in Section 38.10.5):
CREATE FUNCTION execq(text, integer) RETURNS int8
    AS 'filename'
    LANGUAGE C STRICT;

Here is a sample session:

=> SELECT execq('CREATE TABLE a (x integer)', 0);
   execq
   -------
       0
   (1 row)

=> INSERT INTO a VALUES (execq('INSERT INTO a VALUES (0)', 0));
   INSERT 0 1
=> SELECT execq('SELECT * FROM a', 0);
   INFO:  EXECQ:  0    -- inserted by execq
   INFO:  EXECQ:  1    -- returned by execq and inserted by upper
   INSERT
   execq
   -------
       2
   (1 row)

=> SELECT execq('INSERT INTO a SELECT x + 2 FROM a', 1);
   execq
   -------
       1
   (1 row)

=> SELECT execq('SELECT * FROM a', 10);
   INFO:  EXECQ:  0
   INFO:  EXECQ:  1
   INFO:  EXECQ:  2    -- 0 + 2, only one row inserted - as specified
   execq
   -------
       3              -- 10 is the max value only, 3 is the real
       number of rows
   (1 row)

=> DELETE FROM a;
   DELETE 3
=> INSERT INTO a VALUES (execq('SELECT * FROM a', 0) + 1);
   INSERT 0 1
=> SELECT * FROM a;
   x
   ---
   1       -- no rows in a (0) + 1
   (1 row)

=> INSERT INTO a VALUES (execq('SELECT * FROM a', 0) + 1);
   INFO:  EXECQ:  1
   INSERT 0 1
=> SELECT * FROM a;
   x
   ---
-- This demonstrates the data changes visibility rule:

=> INSERT INTO a SELECT execq('SELECT * FROM a', 0) * x FROM a;
INFO: EXECQ: 1
INFO: EXECQ: 2
INFO: EXECQ: 1
INFO: EXECQ: 2
INFO: EXECQ: 2
INSERT 0 2
=> SELECT * FROM a;
x
---
1
2
2 -- 2 rows * 1 (x in first row)
6 -- 3 rows (2 + 1 just inserted) * 2 (x in
second row)
(4 rows)

^~~~~~~
rows visible to execq() in different
invocations
Chapter 48. Background Worker Processes

PostgreSQL can be extended to run user-supplied code in separate processes. Such processes are started, stopped and monitored by `postgres`, which permits them to have a lifetime closely linked to the server’s status. These processes have the option to attach to PostgreSQL’s shared memory area and to connect to databases internally; they can also run multiple transactions serially, just like a regular client-connected server process. Also, by linking to libpq they can connect to the server and behave like a regular client application.

**Warning**

There are considerable robustness and security risks in using background worker processes because, being written in the C language, they have unrestricted access to data. Administrators wishing to enable modules that include background worker process should exercise extreme caution. Only carefully audited modules should be permitted to run background worker processes.

Background workers can be initialized at the time that PostgreSQL is started by including the module name in `shared_preload_libraries`. A module wishing to run a background worker can register it by calling `RegisterBackgroundWorker(BackgroundWorker *worker)` from its `_PG_init()`. Background workers can also be started after the system is up and running by calling the function `RegisterDynamicBackgroundWorker(BackgroundWorker *worker, BackgroundWorkerHandle **handle)` Unlike `RegisterBackgroundWorker`, which can only be called from within the postmaster, `RegisterDynamicBackgroundWorker` must be called from a regular backend or another background worker.

The structure `BackgroundWorker` is defined thus:

```c
typedef void (*bgworker_main_type)(Datum main_arg);
typedef struct BackgroundWorker
{
    char   bgw_name[BGW_MAXLEN];
    char   bgw_type[BGW_MAXLEN];
    int    bgw_flags;
    BgWorkerStartTime bgw_start_time;
    int    bgw_restart_time;  /* in seconds, or BGW_NEVER_RESTART */
    char   bgw_library_name[BGW_MAXLEN];
    char   bgw_function_name[BGW_MAXLEN];
    Datum  bgw_main_arg;
    char   bgw_extra[BGW_EXTRALEN];
    int    bgw_notify_pid;
} BackgroundWorker;
```

`bgw_name` and `bgw_type` are strings to be used in log messages, process listings and similar contexts. `bgw_type` should be the same for all background workers of the same type, so that it is possible to group such workers in a process listing, for example. `bgw_name` on the other hand can contain additional information about the specific process. (Typically, the string for `bgw_name` will contain the type somehow, but that is not strictly required.)

`bgw_flags` is a bitwise-or'd bit mask indicating the capabilities that the module wants. Possible values are:
*Background Worker Processes*

**BGWORKER_SHMEM_ACCESS**

Requests shared memory access. Workers without shared memory access cannot access any of PostgreSQL's shared data structures, such as heavyweight or lightweight locks, shared buffers, or any custom data structures which the worker itself may wish to create and use.

**BGWORKER_BACKEND_DATABASE_CONNECTION**

Requests the ability to establish a database connection through which it can later run transactions and queries. A background worker using `BGWORKER_BACKEND_DATABASE_CONNECTION` to connect to a database must also attach shared memory using `BGWORKER_SHMEM_ACCESS`, or worker start-up will fail.

`bgw_start_time` is the server state during which `postgres` should start the process; it can be one of `BgWorkerStart_PostmasterStart` (start as soon as `postgres` itself has finished its own initialization; processes requesting this are not eligible for database connections), `BgWorkerStart_ConsistentState` (start as soon as a consistent state has been reached in a hot standby, allowing processes to connect to databases and run read-only queries), and `BgWorkerStart_RecoveryFinished` (start as soon as the system has entered normal read-write state). Note the last two values are equivalent in a server that's not a hot standby. Note that this setting only indicates when the processes are to be started; they do not stop when a different state is reached.

`bgw_restart_time` is the interval, in seconds, that `postgres` should wait before restarting the process, in case it crashes. It can be any positive value, or `BGW_NEVER_RESTART`, indicating not to restart the process in case of a crash.

`bgw_library_name` is the name of a library in which the initial entry point for the background worker should be sought. The named library will be dynamically loaded by the worker process and `bgw_function_name` will be used to identify the function to be called. If loading a function from the core code, this must be set to "postgres".

`bgw_function_name` is the name of a function in a dynamically loaded library which should be used as the initial entry point for a new background worker.

`bgw_main_arg` is the Datum argument to the background worker main function. This main function should take a single argument of type Datum and return void. `bgw_main_arg` will be passed as the argument. In addition, the global variable `MyBgworkerEntry` points to a copy of the `BackgroundWorker` structure passed at registration time; the worker may find it helpful to examine this structure.

On Windows (and anywhere else where `EXEC_BACKEND` is defined) or in dynamic background workers it is not safe to pass a Datum by reference, only by value. If an argument is required, it is safest to pass an int32 or other small value and use that as an index into an array allocated in shared memory. If a value like a cstring or text is passed then the pointer won't be valid from the new background worker process.

`bgw_extra` can contain extra data to be passed to the background worker. Unlike `bgw_main_arg`, this data is not passed as an argument to the worker's main function, but it can be accessed via `MyBgworkerEntry`, as discussed above.

`bgw_notify_pid` is the PID of a PostgreSQL backend process to which the postmaster should send SIGUSR1 when the process is started or exits. It should be 0 for workers registered at postmaster startup time, or when the backend registering the worker does not wish to wait for the worker to start up. Otherwise, it should be initialized to `MyProcPid`.

Once running, the process can connect to a database by calling `BackgroundWorkerInitializeConnection(char *dbname, char *username, uint32 flags)` or `BackgroundWorkerInitializeConnectionByOid(Oid dboid, Oid useroid, uint32 flags)`. This allows the process to run transactions and queries using the SPI interface. If `dbname` is NULL or `dboid` is `InvalidOid`, the session is not connected to any particular database, but
Background Worker Processes

shared catalogs can be accessed. If $\text{username}$ is NULL or $\text{useroid}$ is InvalidOid, the process will run as the superuser created during initdb. If BGWORKER_BYPASS_ALLOWCONN is specified as flags it is possible to bypass the restriction to connect to databases not allowing user connections. A background worker can only call one of these two functions, and only once. It is not possible to switch databases.

Signals are initially blocked when control reaches the background worker’s main function, and must be unblocked by it; this is to allow the process to customize its signal handlers, if necessary. Signals can be unblocked in the new process by calling BackgroundWorkerUnblockSignals and blocked by calling BackgroundWorkerBlockSignals.

If $\text{bgw\_restart\_time}$ for a background worker is configured as BGW_NEVER_RESTART, or if it exits with an exit code of 0 or is terminated by TerminateBackgroundWorker, it will be automatically unregistered by the postmaster on exit. Otherwise, it will be restarted after the time period configured via $\text{bgw\_restart\_time}$, or immediately if the postmaster reinitializes the cluster due to a backend failure. Backends which need to suspend execution only temporarily should use an interruptible sleep rather than exiting; this can be achieved by calling WaitLatch(). Make sure the WL_POSTMASTER_DEATH flag is set when calling that function, and verify the return code for a prompt exit in the emergency case that postgres itself has terminated.

When a background worker is registered using the RegisterDynamicBackgroundWorker function, it is possible for the backend performing the registration to obtain information regarding the status of the worker. Backends wishing to do this should pass the address of a BackgroundWorkerHandle * as the second argument to RegisterDynamicBackgroundWorker. If the worker is successfully registered, this pointer will be initialized with an opaque handle that can subsequently be passed to GetBackgroundWorkerPid(BackgroundWorkerHandle *, pid_t *) or TerminateBackgroundWorker(BackgroundWorkerHandle *). GetBackgroundWorkerPid can be used to poll the status of the worker: a return value of BGWH_NOT_YET_STARTED indicates that the worker has not yet been started by the postmaster; BGWH_STOPPED indicates that it has been started but is no longer running; and BGWH_STARTED indicates that it is currently running. In this last case, the PID will also be returned via the second argument. TerminateBackgroundWorker causes the postmaster to send SIGTERM to the worker if it is running, and to unregister it as soon as it is not.

In some cases, a process which registers a background worker may wish to wait for the worker to start up. This can be accomplished by initializing bgw_notify_pid to MyProcPid and then passing the BackgroundWorkerHandle * obtained at registration time to WaitForBackgroundWorkerStartup(BackgroundWorkerHandle *handle, pid_t *) function. This function will block until the postmaster has attempted to start the background worker, or until the postmaster dies. If the background worker is running, the return value will be BGWH_STARTED, and the PID will be written to the provided address. Otherwise, the return value will be BGWH_STOPPED or BGWH_POSTMASTER_DIED.

A process can also wait for a background worker to shut down, by using the WaitForBackgroundWorkerShutdown(BackgroundWorkerHandle *handle) function and passing the BackgroundWorkerHandle * obtained at registration. This function will block until the background worker exits, or postmaster dies. When the background worker exits, the return value is BGWH_STOPPED, if postmaster dies it will return BGWH_POSTMASTER_DIED.

If a background worker sends asynchronous notifications with the NOTIFY command via the Server Programming Interface (SPI), it should call ProcessCompletedNotifies explicitly after committing the enclosing transaction so that any notifications can be delivered. If a background worker registers to receive asynchronous notifications with the LISTEN through SPI, the worker will log those notifications, but there is no programmatic way for the worker to intercept and respond to those notifications.

The $\text{src/test/modules/worker_spi}$ module contains a working example, which demonstrates some useful techniques.

The maximum number of registered background workers is limited by max_worker_processes.
Chapter 49. Logical Decoding

PostgreSQL provides infrastructure to stream the modifications performed via SQL to external consumers. This functionality can be used for a variety of purposes, including replication solutions and auditing.

Changes are sent out in streams identified by logical replication slots.

The format in which those changes are streamed is determined by the output plugin used. An example plugin is provided in the PostgreSQL distribution. Additional plugins can be written to extend the choice of available formats without modifying any core code. Every output plugin has access to each individual new row produced by \texttt{INSERT} and the new row version created by \texttt{UPDATE}. Availability of old row versions for \texttt{UPDATE} and \texttt{DELETE} depends on the configured replica identity (see \texttt{REPLICA IDENTITY}).

Changes can be consumed either using the streaming replication protocol (see Section 53.4 and Section 49.3), or by calling functions via SQL (see Section 49.4). It is also possible to write additional methods of consuming the output of a replication slot without modifying core code (see Section 49.7).

49.1. Logical Decoding Examples

The following example demonstrates controlling logical decoding using the SQL interface.

Before you can use logical decoding, you must set wal_level to logical and max_replication_slots to at least 1. Then, you should connect to the target database (in the example below, postgres) as a superuser.

```sql
postgres=# -- Create a slot named 'regression_slot' using the output plugin 'test_decoding'
postgres=# SELECT * FROM pg_create_logical_replication_slot('regression_slot', 'test_decoding');

+-------------------+-----------+
| slot_name | lsn       |
+-------------------+-----------+
| regression_slot   | 0/16B1970 |
+-------------------+-----------+

postgres=# SELECT slot_name, plugin, slot_type, database, active, restart_lsn, confirmed_flush_lsn FROM pg_replication_slots;

+-------------------+---------------+-----------+----------+--------+-------------------+-------------------+
| slot_name | plugin     | slot_type | database | active | restart_lsn | confirmed_flush_lsn |
+-------------------+---------------+-----------+----------+--------+-------------------+-------------------+
| regression_slot   | test_decoding | logical   | postgres | f      | 0/16A4408 | 0/16A4440           |
+-------------------+---------------+-----------+----------+--------+-------------------+-------------------+

postgres=# -- There are no changes to see yet
postgres=# SELECT * FROM pg_logical_slot_get_changes('regression_slot', NULL, NULL);

lsn | xid | data
-----+-----+------
(0 rows)

postgres=# CREATE TABLE data(id serial primary key, data text);
CREATE TABLE
```
Logical Decoding

```
postgres=# -- DDL isn't replicated, so all you'll see is the transaction
postgres=# SELECT * FROM
pg_logical_slot_get_changes('regression_slot', NULL, NULL);

| lsn     | xid  |     data         |
|---------+------+------------------|
| 0/BA2DA58 | 10297 | BEGIN 10297      |
| 0/BA5A5A0 | 10297 | COMMIT 10297     |

(2 rows)

postgres=# -- Once changes are read, they're consumed and not emitted
postgres=# -- in a subsequent call:
postgres=# SELECT * FROM
pg_logical_slot_get_changes('regression_slot', NULL, NULL);

lsn | xid | data
-----+-----+------
(0 rows)

postgres=# BEGIN;
postgres=# INSERT INTO data(data) VALUES('1');
postgres=# INSERT INTO data(data) VALUES('2');
postgres=# COMMIT;

postgres=# SELECT * FROM
pg_logical_slot_get_changes('regression_slot', NULL, NULL);

| lsn     | xid  |                          data                          |
|---------+------+----------------------------------------------------------------|
| 0/BA5A688 | 10298 | BEGIN 10298                                                            |
| 0/BA5A6F0 | 10298 | table public.data: INSERT: id[integer]:1 data[text]:'1'                |
| 0/BA5A7F8 | 10298 | table public.data: INSERT: id[integer]:2 data[text]:'2'                |
| 0/BA5A8A8 | 10298 | COMMIT 10298                                                            |

(4 rows)

postgres=# INSERT INTO data(data) VALUES('3');

postgres=# -- You can also peek ahead in the change stream without consuming changes
postgres=# SELECT * FROM
pg_logical_slot_peek_changes('regression_slot', NULL, NULL);

| lsn     | xid  |                          data                          |
|---------+------+----------------------------------------------------------------|
| 0/BA5A8E0 | 10299 | BEGIN 10299                                                            |
| 0/BA5A8E0 | 10299 | table public.data: INSERT: id[integer]:3 data[text]:'3'                |
| 0/BA5A990 | 10299 | COMMIT 10299                                                            |

(3 rows)

postgres=# -- The next call to pg_logical_slot_peek_changes() returns the same changes again
postgres=# SELECT * FROM
pg_logical_slot_peek_changes('regression_slot', NULL, NULL);

```

1333
49.2. Logical Decoding Concepts

49.2.1. Logical Decoding
Logical decoding is the process of extracting all persistent changes to a database's tables into a coherent, easy to understand format which can be interpreted without detailed knowledge of the database's internal state.

In PostgreSQL, logical decoding is implemented by decoding the contents of the write-ahead log, which describe changes on a storage level, into an application-specific form such as a stream of tuples or SQL statements.

### 49.2.2. Replication Slots

In the context of logical replication, a slot represents a stream of changes that can be replayed to a client in the order they were made on the origin server. Each slot streams a sequence of changes from a single database.

**Note**

PostgreSQL also has streaming replication slots (see Section 26.2.5), but they are used somewhat differently there.

A replication slot has an identifier that is unique across all databases in a PostgreSQL cluster. Slots persist independently of the connection using them and are crash-safe.

A logical slot will emit each change just once in normal operation. The current position of each slot is persisted only at checkpoint, so in the case of a crash the slot may return to an earlier LSN, which will then cause recent changes to be resent when the server restarts. Logical decoding clients are responsible for avoiding ill effects from handling the same message more than once. Clients may wish to record the last LSN they saw when decoding and skip over any repeated data or (when using the replication protocol) request that decoding start from that LSN rather than letting the server determine the start point. The Replication Progress Tracking feature is designed for this purpose, refer to replication origins.

Multiple independent slots may exist for a single database. Each slot has its own state, allowing different consumers to receive changes from different points in the database change stream. For most applications, a separate slot will be required for each consumer.

A logical replication slot knows nothing about the state of the receiver(s). It's even possible to have multiple different receivers using the same slot at different times; they'll just get the changes following on from when the last receiver stopped consuming them. Only one receiver may consume changes from a slot at any given time.

**Caution**

Replication slots persist across crashes and know nothing about the state of their consumer(s). They will prevent removal of required resources even when there is no connection using them. This consumes storage because neither required WAL nor required rows from the system catalogs can be removed by VACUUM as long as they are required by a replication slot. In extreme cases this could cause the database to shut down to prevent transaction ID wraparound (see Section 24.1.5). So if a slot is no longer required it should be dropped.

### 49.2.3. Output Plugins

Output plugins transform the data from the write-ahead log's internal representation into the format the consumer of a replication slot desires.
49.2.4. Exported Snapshots

When a new replication slot is created using the streaming replication interface (see CREATE_REPLICATION_SLOT), a snapshot is exported (see Section 9.26.5), which will show exactly the state of the database after which all changes will be included in the change stream. This can be used to create a new replica by using SET TRANSACTION_SNAPSHOT to read the state of the database at the moment the slot was created. This transaction can then be used to dump the database's state at that point in time, which afterwards can be updated using the slot's contents without losing any changes.

Creation of a snapshot is not always possible. In particular, it will fail when connected to a hot standby. Applications that do not require snapshot export may suppress it with the NOEXPORT_SNAPSHOT option.

49.3. Streaming Replication Protocol Interface

The commands

- CREATE_REPLICATION_SLOT slot_name LOGICAL output_plugin
- DROP_REPLICATION_SLOT slot_name [ WAIT ]
- START_REPLICATION SLOT slot_name LOGICAL ...

are used to create, drop, and stream changes from a replication slot, respectively. These commands are only available over a replication connection; they cannot be used via SQL. See Section 53.4 for details on these commands.

The command pg_recvlogical can be used to control logical decoding over a streaming replication connection. (It uses these commands internally.)

49.4. Logical Decoding SQL Interface

See Section 9.26.6 for detailed documentation on the SQL-level API for interacting with logical decoding.

Synchronous replication (see Section 26.2.8) is only supported on replication slots used over the streaming replication interface. The function interface and additional, non-core interfaces do not support synchronous replication.

49.5. System Catalogs Related to Logical Decoding

The pg_replication_slots view and the pg_stat_replication view provide information about the current state of replication slots and streaming replication connections respectively. These views apply to both physical and logical replication.

49.6. Logical Decoding Output Plugins

An example output plugin can be found in the contrib/test_decoding subdirectory of the PostgreSQL source tree.

49.6.1. Initialization Function
An output plugin is loaded by dynamically loading a shared library with the output plugin’s name as the library base name. The normal library search path is used to locate the library. To provide the required output plugin callbacks and to indicate that the library is actually an output plugin it needs to provide a function named \_PG_output_plugin_init. This function is passed a struct that needs to be filled with the callback function pointers for individual actions.

```c
typedef struct OutputPluginCallbacks {
    LogicalDecodeStartupCB startup_cb;
    LogicalDecodeBeginCB begin_cb;
    LogicalDecodeChangeCB change_cb;
    LogicalDecodeTruncateCB truncate_cb;
    LogicalDecodeCommitCB commit_cb;
    LogicalDecodeMessageCB message_cb;
    LogicalDecodeFilterByOriginCB filter_by_origin_cb;
    LogicalDecodeShutdownCB shutdown_cb;
} OutputPluginCallbacks;

typedef void (*LogicalOutputPluginInit) (struct OutputPluginCallbacks *cb);
```

The `begin_cb`, `change_cb` and `commit_cb` callbacks are required, while `startup_cb`, `filter_by_origin_cb`, `truncate_cb`, and `shutdown_cb` are optional. If `truncate_cb` is not set but a TRUNCATE is to be decoded, the action will be ignored.

### 49.6.2. Capabilities

To decode, format and output changes, output plugins can use most of the backend’s normal infrastructure, including calling output functions. Read only access to relations is permitted as long as only relations are accessed that either have been created by `initdb` in the `pg_catalog` schema, or have been marked as user provided catalog tables using

```
ALTER TABLE user_catalog_table SET (user_catalog_table = true);
CREATE TABLE another_catalog_table(data text) WITH (user_catalog_table = true);
```

Any actions leading to transaction ID assignment are prohibited. That, among others, includes writing to tables, performing DDL changes, and calling `txid_current()`.

### 49.6.3. Output Modes

Output plugin callbacks can pass data to the consumer in nearly arbitrary formats. For some use cases, like viewing the changes via SQL, returning data in a data type that can contain arbitrary data (e.g., `bytea`) is cumbersome. If the output plugin only outputs textual data in the server’s encoding, it can declare that by setting `OutputPluginOptions.output_type` to `OUTPUT_PLUGIN_TEXTUAL_OUTPUT` instead of `OUTPUT_PLUGIN_BINARY_OUTPUT` in the startup callback. In that case, all the data has to be in the server’s encoding so that a `text` datum can contain it. This is checked in assertion-enabled builds.

### 49.6.4. Output Plugin Callbacks

An output plugin gets notified about changes that are happening via various callbacks it needs to provide.

Concurrent transactions are decoded in commit order, and only changes belonging to a specific transaction are decoded between the `begin` and `commit` callbacks. Transactions that were rolled back...
explicitly or implicitly never get decoded. Successful savepoints are folded into the transaction containing them in the order they were executed within that transaction.

**Note**

Only transactions that have already safely been flushed to disk will be decoded. That can lead to a COMMIT not immediately being decoded in a directly following `pg_logical_slot_get_changes()` when synchronous_commit is set to off.

### 49.6.4.1. Startup Callback

The optional startup_cb callback is called whenever a replication slot is created or asked to stream changes, independent of the number of changes that are ready to be put out.

```c
typedef void (*LogicalDecodeStartupCB) (struct LogicalDecodingContext *ctx,
                                        OutputPluginOptions *options,
                                        bool is_init);
```

The is_init parameter will be true when the replication slot is being created and false otherwise. options points to a struct of options that output plugins can set:

```c
typedef struct OutputPluginOptions {
    OutputPluginOutputType output_type;
    bool receive_rewrites;
} OutputPluginOptions;
```

output_type has to either be set to OUTPUT_PLUGIN_TEXTUAL_OUTPUT or OUTPUT_PLUGIN_BINARY_OUTPUT. See also Section 49.6.3. If receive_rewrites is true, the output plugin will also be called for changes made by heap rewrites during certain DDL operations. These are of interest to plugins that handle DDL replication, but they require special handling.

The startup callback should validate the options present in `ctx->output_plugin_options`. If the output plugin needs to have a state, it can use `ctx->output_plugin_private` to store it.

### 49.6.4.2. Shutdown Callback

The optional shutdown_cb callback is called whenever a formerly active replication slot is not used anymore and can be used to deallocate resources private to the output plugin. The slot isn't necessarily being dropped, streaming is just being stopped.

```c
typedef void (*LogicalDecodeShutdownCB) (struct LogicalDecodingContext *ctx);
```

### 49.6.4.3. Transaction Begin Callback

The required begin_cb callback is called whenever a start of a committed transaction has been decoded. Aborted transactions and their contents never get decoded.

```c
typedef void (*LogicalDecodeBeginCB) (struct LogicalDecodingContext *ctx,
                                      1338
```
The `txn` parameter contains meta information about the transaction, like the time stamp at which it has been committed and its XID.

### 49.6.4.4. Transaction End Callback

The required `commit_cb` callback is called whenever a transaction commit has been decoded. The `change_cb` callbacks for all modified rows will have been called before this, if there have been any modified rows.

```c
typedef void (*LogicalDecodeCommitCB) (struct LogicalDecodingContext *ctx,
                                       ReorderBufferTXN *txn,
                                       XLogRecPtr commit_lsn);
```

### 49.6.4.5. Change Callback

The required `change_cb` callback is called for every individual row modification inside a transaction, may it be an `INSERT`, `UPDATE`, or `DELETE`. Even if the original command modified several rows at once the callback will be called individually for each row.

```c
typedef void (*LogicalDecodeChangeCB) (struct LogicalDecodingContext *ctx,
                                        ReorderBufferTXN *txn,
                                        Relation relation,
                                        ReorderBufferChange *change);
```

The `ctx` and `txn` parameters have the same contents as for the `begin_cb` and `commit_cb` callbacks, but additionally the relation descriptor `relation` points to the relation the row belongs to and a struct `change` describing the row modification are passed in.

**Note**

Only changes in user defined tables that are not unlogged (see UNLOGGED) and not temporary (see TEMPORARY or TEMP) can be extracted using logical decoding.

### 49.6.4.6. Truncate Callback

The `truncate_cb` callback is called for a TRUNCATE command.

```c
typedef void (*LogicalDecodeTruncateCB) (struct LogicalDecodingContext *ctx,
                                          ReorderBufferTXN *txn,
                                          int nrelations,
                                          Relation relations[],
                                          ReorderBufferChange *change);
```

The parameters are analogous to the `change_cb` callback. However, because TRUNCATE actions on tables connected by foreign keys need to be executed together, this callback receives an array of relations instead of just a single one. See the description of the TRUNCATE statement for details.
49.6.4.7. Origin Filter Callback

The optional `filter_by_origin_cb` callback is called to determine whether data that has been replayed from `origin_id` is of interest to the output plugin.

```c
typedef bool (*LogicalDecodeFilterByOriginCB) (struct LogicalDecodingContext *ctx, RepOriginId origin_id);
```

The `ctx` parameter has the same contents as for the other callbacks. No information but the origin is available. To signal that changes originating on the passed in node are irrelevant, return true, causing them to be filtered away; false otherwise. The other callbacks will not be called for transactions and changes that have been filtered away.

This is useful when implementing cascading or multidirectional replication solutions. Filtering by the origin allows to prevent replicating the same changes back and forth in such setups. While transactions and changes also carry information about the origin, filtering via this callback is noticeably more efficient.

49.6.4.8. Generic Message Callback

The optional `message_cb` callback is called whenever a logical decoding message has been decoded.

```c
typedef void (*LogicalDecodeMessageCB) (struct LogicalDecodingContext *ctx, ReorderBufferTXN *txn, XLogRecPtr message_lsn, bool transactional, const char *prefix, Size message_size, const char *message);
```

The `txn` parameter contains meta information about the transaction, like the time stamp at which it has been committed and its XID. Note however that it can be NULL when the message is non-transactional and the XID was not assigned yet in the transaction which logged the message. The `lsn` has WAL location of the message. The `transactional` says if the message was sent as transactional or not. The `prefix` is arbitrary null-terminated prefix which can be used for identifying interesting messages for the current plugin. And finally the `message` parameter holds the actual message of `message_size` size.

Extra care should be taken to ensure that the prefix the output plugin considers interesting is unique. Using name of the extension or the output plugin itself is often a good choice.

49.6.5. Functions for Producing Output

To actually produce output, output plugins can write data to the `StringInfo` output buffer in `ctx->out` when inside the `begin_cb`, `commit_cb`, or `change_cb` callbacks. Before writing to the output buffer, `OutputPluginPrepareWrite(ctx, last_write)` has to be called, and after finishing writing to the buffer, `OutputPluginWrite(ctx, last_write)` has to be called to perform the write. The `last_write` indicates whether a particular write was the callback’s last write.

The following example shows how to output data to the consumer of an output plugin:

```c
OutputPluginPrepareWrite(ctx, true);
appendStringInfo(ctx->out, "BEGIN %u", txn->xid);
```
49.7. Logical Decoding Output Writers

It is possible to add more output methods for logical decoding. For details, see src/backend/replication/logical/logicalfuncs.c. Essentially, three functions need to be provided: one to read WAL, one to prepare writing output, and one to write the output (see Section 49.6.5).

49.8. Synchronous Replication Support for Logical Decoding

Logical decoding can be used to build synchronous replication solutions with the same user interface as synchronous replication for streaming replication. To do this, the streaming replication interface (see Section 49.3) must be used to stream out data. Clients have to send Standby status update (F) (see Section 53.4) messages, just like streaming replication clients do.

Note

A synchronous replica receiving changes via logical decoding will work in the scope of a single database. Since, in contrast to that, synchronous_standby_names currently is server wide, this means this technique will not work properly if more than one database is actively used.
Chapter 50. Replication Progress Tracking

Replication origins are intended to make it easier to implement logical replication solutions on top of logical decoding. They provide a solution to two common problems:

- How to safely keep track of replication progress
- How to change replication behavior based on the origin of a row; for example, to prevent loops in bi-directional replication setups

Replication origins have just two properties, a name and an OID. The name, which is what should be used to refer to the origin across systems, is free-form text. It should be used in a way that makes conflicts between replication origins created by different replication solutions unlikely; e.g. by prefixing the replication solution's name to it. The OID is used only to avoid having to store the long version in situations where space efficiency is important. It should never be shared across systems.

Replication origins can be created using the function `pg_replication_origin_create();` dropped using `pg_replication_origin_drop();` and seen in the `pg_replication_origin` system catalog.

One nontrivial part of building a replication solution is to keep track of replay progress in a safe manner. When the applying process, or the whole cluster, dies, it needs to be possible to find out up to where data has successfully been replicated. Naive solutions to this, such as updating a row in a table for every replayed transaction, have problems like run-time overhead and database bloat.

Using the replication origin infrastructure a session can be marked as replaying from a remote node (using the `pg_replication_origin_session_setup()` function). Additionally the LSN and commit time stamp of every source transaction can be configured on a per transaction basis using `pg_replication_origin_xact_setup()`. If that's done replication progress will persist in a crash safe manner. Replay progress for all replication origins can be seen in the `pg_replication_origin_status` view. An individual origin's progress, e.g. when resuming replication, can be acquired using `pg_replication_origin_progress()` for any origin or `pg_replication_origin_session_progress()` for the origin configured in the current session.

In replication topologies more complex than replication from exactly one system to one other system, another problem can be that it is hard to avoid replicating replayed rows again. That can lead both to cycles in the replication and inefficiencies. Replication origins provide an optional mechanism to recognize and prevent that. When configured using the functions referenced in the previous paragraph, every change and transaction passed to output plugin callbacks (see Section 49.6) generated by the session is tagged with the replication origin of the generating session. This allows treating them differently in the output plugin, e.g. ignoring all but locally-originating rows. Additionally the `filter_by_origin_cb` callback can be used to filter the logical decoding change stream based on the source. While less flexible, filtering via that callback is considerably more efficient than doing it in the output plugin.
Part VI. Reference

The entries in this Reference are meant to provide in reasonable length an authoritative, complete, and formal summary about their respective subjects. More information about the use of PostgreSQL, in narrative, tutorial, or example form, can be found in other parts of this book. See the cross-references listed on each reference page.

The reference entries are also available as traditional “man” pages.
# Table of Contents

I. SQL Commands ................................................................. 1348  
  ABORT ................................................................. 1352  
  ALTER AGGREGATE ................................................................. 1353  
  ALTER COLLATION ................................................................. 1355  
  ALTER CONVERSION ................................................................. 1357  
  ALTER DATABASE ................................................................. 1359  
  ALTER DEFAULT PRIVILEGES ................................................................. 1362  
  ALTER DOMAIN ................................................................. 1365  
  ALTER EVENT TRIGGER ......................................................... 1368  
  ALTER EXTENSION ................................................................. 1369  
  ALTER FOREIGN DATA WRAPPER ................................................................. 1373  
  ALTER FOREIGN TABLE ................................................................. 1375  
  ALTER FUNCTION ................................................................. 1380  
  ALTER GROUP ................................................................. 1384  
  ALTER INDEX ................................................................. 1386  
  ALTER LANGUAGE ................................................................. 1389  
  ALTER LARGE OBJECT ................................................................. 1390  
  ALTER MATERIALIZED VIEW ................................................................. 1391  
  ALTER OPERATOR ................................................................. 1393  
  ALTER OPERATOR CLASS ................................................................. 1395  
  ALTER OPERATOR FAMILY ................................................................. 1396  
  ALTER POLICY ................................................................. 1400  
  ALTER PROCEDURE ................................................................. 1402  
  ALTER PUBLICATION ................................................................. 1405  
  ALTER ROLE ................................................................. 1407  
  ALTER ROUTINE ................................................................. 1411  
  ALTER RULE ................................................................. 1413  
  ALTER SCHEMA ................................................................. 1414  
  ALTER SEQUENCE ................................................................. 1415  
  ALTER SERVER ................................................................. 1418  
  ALTER STATISTICS ................................................................. 1420  
  ALTER SUBSCRIPTION ................................................................. 1421  
  ALTER SYSTEM ................................................................. 1423  
  ALTER TABLE ................................................................. 1425  
  ALTER TABLESPACE ................................................................. 1441  
  ALTER TEXT SEARCH CONFIGURATION ................................................................. 1443  
  ALTER TEXT SEARCH DICTIONARY ................................................................. 1445  
  ALTER TEXT SEARCH PARSER ................................................................. 1447  
  ALTER TEXT SEARCH TEMPLATE ................................................................. 1448  
  ALTER TRIGGER ................................................................. 1449  
  ALTER TYPE ................................................................. 1451  
  ALTER USER ................................................................. 1455  
  ALTER USER MAPPING ................................................................. 1456  
  ALTER VIEW ................................................................. 1457  
  ANALYZE ................................................................. 1459  
  BEGIN ................................................................. 1462  
  CALL ................................................................. 1464  
  CHECKPOINT ................................................................. 1465  
  CLOSE ................................................................. 1466  
  CLUSTER ................................................................. 1467  
  COMMENT ................................................................. 1469  
  COMMIT ................................................................. 1474  
  COMMIT PREPARED ................................................................. 1475  
  COPY ................................................................. 1476  
  CREATE ACCESS METHOD ................................................................. 1486
<table>
<thead>
<tr>
<th>Command</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>CREATE AGGREGATE</td>
<td>1487</td>
</tr>
<tr>
<td>CREATE CAST</td>
<td>1495</td>
</tr>
<tr>
<td>CREATE COLLATION</td>
<td>1499</td>
</tr>
<tr>
<td>CREATE CONVERSION</td>
<td>1501</td>
</tr>
<tr>
<td>CREATE DATABASE</td>
<td>1503</td>
</tr>
<tr>
<td>CREATE DOMAIN</td>
<td>1506</td>
</tr>
<tr>
<td>CREATE EVENT TRIGGER</td>
<td>1509</td>
</tr>
<tr>
<td>CREATE EXTENSION</td>
<td>1511</td>
</tr>
<tr>
<td>CREATE FOREIGN DATA WRAPPER</td>
<td>1513</td>
</tr>
<tr>
<td>CREATE FOREIGN TABLE</td>
<td>1515</td>
</tr>
<tr>
<td>CREATE FUNCTION</td>
<td>1519</td>
</tr>
<tr>
<td>CREATE GROUP</td>
<td>1527</td>
</tr>
<tr>
<td>CREATE INDEX</td>
<td>1528</td>
</tr>
<tr>
<td>CREATE LANGUAGE</td>
<td>1536</td>
</tr>
<tr>
<td>CREATE MATERIALIZED VIEW</td>
<td>1539</td>
</tr>
<tr>
<td>CREATE OPERATOR</td>
<td>1541</td>
</tr>
<tr>
<td>CREATE OPERATOR CLASS</td>
<td>1544</td>
</tr>
<tr>
<td>CREATE OPERATOR FAMILY</td>
<td>1547</td>
</tr>
<tr>
<td>CREATE POLICY</td>
<td>1548</td>
</tr>
<tr>
<td>CREATE PROCEDURE</td>
<td>1554</td>
</tr>
<tr>
<td>CREATE PUBLICATION</td>
<td>1557</td>
</tr>
<tr>
<td>CREATE ROLE</td>
<td>1559</td>
</tr>
<tr>
<td>CREATE RULE</td>
<td>1564</td>
</tr>
<tr>
<td>CREATE SCHEMA</td>
<td>1567</td>
</tr>
<tr>
<td>CREATE SEQUENCE</td>
<td>1570</td>
</tr>
<tr>
<td>CREATE SERVER</td>
<td>1574</td>
</tr>
<tr>
<td>CREATE STATISTICS</td>
<td>1576</td>
</tr>
<tr>
<td>CREATE SUBSCRIPTION</td>
<td>1578</td>
</tr>
<tr>
<td>CREATE TABLE</td>
<td>1581</td>
</tr>
<tr>
<td>CREATE TABLE AS</td>
<td>1601</td>
</tr>
<tr>
<td>CREATE TABLESPACE</td>
<td>1604</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH CONFIGURATION</td>
<td>1606</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH DICTIONARY</td>
<td>1607</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH PARSER</td>
<td>1609</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TEMPLATE</td>
<td>1611</td>
</tr>
<tr>
<td>CREATE TRANSFORM</td>
<td>1612</td>
</tr>
<tr>
<td>CREATE TRIGGER</td>
<td>1614</td>
</tr>
<tr>
<td>CREATE TYPE</td>
<td>1621</td>
</tr>
<tr>
<td>CREATE USER</td>
<td>1630</td>
</tr>
<tr>
<td>CREATE USER MAPPING</td>
<td>1631</td>
</tr>
<tr>
<td>CREATE VIEW</td>
<td>1633</td>
</tr>
<tr>
<td>DEALLOCATE</td>
<td>1638</td>
</tr>
<tr>
<td>DECLARE</td>
<td>1639</td>
</tr>
<tr>
<td>DELETE</td>
<td>1642</td>
</tr>
<tr>
<td>DISCARD</td>
<td>1645</td>
</tr>
<tr>
<td>DO</td>
<td>1646</td>
</tr>
<tr>
<td>DROP ACCESS METHOD</td>
<td>1648</td>
</tr>
<tr>
<td>DROP AGGREGATE</td>
<td>1649</td>
</tr>
<tr>
<td>DROP CAST</td>
<td>1651</td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>1652</td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>1653</td>
</tr>
<tr>
<td>DROP DATABASE</td>
<td>1654</td>
</tr>
<tr>
<td>DROP DOMAIN</td>
<td>1655</td>
</tr>
<tr>
<td>DROP EVENT TRIGGER</td>
<td>1656</td>
</tr>
<tr>
<td>DROP EXTENSION</td>
<td>1657</td>
</tr>
<tr>
<td>DROP FOREIGN DATA WRAPPER</td>
<td>1658</td>
</tr>
<tr>
<td>DROP FOREIGN TABLE</td>
<td>1659</td>
</tr>
<tr>
<td>DROP FUNCTION</td>
<td>1660</td>
</tr>
</tbody>
</table>
DROP GROUP .............................................................. 1662
DROP INDEX ............................................................ 1663
DROP LANGUAGE ....................................................... 1665
DROP MATERIALIZED VIEW ........................................... 1666
DROP OPERATOR ......................................................... 1667
DROP OPERATOR CLASS ............................................... 1669
DROP OPERATOR FAMILY ............................................ 1671
DROP OWNED ............................................................ 1673
DROP POLICY ............................................................ 1674
DROP PROCEDURE ....................................................... 1675
DROP PUBLICATION .................................................... 1677
DROP ROLE ............................................................... 1678
DROP ROUTINE .......................................................... 1679
DROP RULE ............................................................... 1680
DROP SCHEMA ............................................................ 1681
DROP SEQUENCE ........................................................ 1682
DROP SERVER ........................................................... 1683
DROP STATISTICS ....................................................... 1684
DROP SUBSCRIPTION ................................................... 1685
DROP TABLE ............................................................. 1687
DROP TABLESPACE ...................................................... 1688
DROP TEXT SEARCH CONFIGURATION ............................... 1689
DROP TEXT SEARCH DICTIONARY .................................... 1690
DROP TEXT SEARCH PARSER .......................................... 1691
DROP TEXT SEARCH TEMPLATE ....................................... 1692
DROP TRANSFORM ....................................................... 1693
DROP TRIGGER ........................................................... 1694
DROP TYPE ............................................................... 1695
DROP USER ............................................................... 1696
DROP USER MAPPING .................................................. 1697
DROP VIEW ............................................................... 1698
END ............................................................................. 1699
EXECUTE .................................................................... 1700
EXPLAIN ...................................................................... 1701
FETCH ........................................................................ 1706
GRANT ........................................................................ 1710
IMPORT FOREIGN SCHEMA ............................................... 1717
INSERT ...................................................................... 1719
LISTEN ........................................................................ 1726
LOAD .......................................................................... 1728
LOCK .......................................................................... 1729
MOVE ........................................................................ 1732
NOTIFY ........................................................................ 1734
PREPARE ..................................................................... 1737
PREPARE TRANSACTION ................................................ 1740
REASSIGN OWNED ....................................................... 1742
REFRESH MATERIALIZED VIEW ....................................... 1743
REINDEX ..................................................................... 1745
RELEASE SAVEPOINT .................................................... 1748
RESET .......................................................................... 1750
REVOKE ...................................................................... 1751
ROLLBACK .................................................................... 1755
ROLLBACK PREPARED ................................................... 1756
ROLLBACK TO SAVEPOINT ............................................. 1757
SAVEPOINT ................................................................... 1759
SECURITY LABEL .......................................................... 1761
SELECT ........................................................................ 1764
SELECT INTO ............................................................... 1785
SQL Commands

This part contains reference information for the SQL commands supported by PostgreSQL. By “SQL” the language in general is meant; information about the standards conformance and compatibility of each command can be found on the respective reference page.

Table of Contents

<table>
<thead>
<tr>
<th>Command</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>ABORT</td>
<td>1352</td>
</tr>
<tr>
<td>ALTER AGGREGATE</td>
<td>1353</td>
</tr>
<tr>
<td>ALTER COLLATION</td>
<td>1355</td>
</tr>
<tr>
<td>ALTER CONVERSION</td>
<td>1357</td>
</tr>
<tr>
<td>ALTER DATABASE</td>
<td>1359</td>
</tr>
<tr>
<td>ALTER DEFAULT PRIVILEGES</td>
<td>1362</td>
</tr>
<tr>
<td>ALTER DOMAIN</td>
<td>1365</td>
</tr>
<tr>
<td>ALTER EVENT TRIGGER</td>
<td>1368</td>
</tr>
<tr>
<td>ALTER EXTENSION</td>
<td>1369</td>
</tr>
<tr>
<td>ALTER FOREIGN DATA WRAPPER</td>
<td>1373</td>
</tr>
<tr>
<td>ALTER FOREIGN TABLE</td>
<td>1375</td>
</tr>
<tr>
<td>ALTER FUNCTION</td>
<td>1380</td>
</tr>
<tr>
<td>ALTER GROUP</td>
<td>1384</td>
</tr>
<tr>
<td>ALTER INDEX</td>
<td>1386</td>
</tr>
<tr>
<td>ALTER LANGUAGE</td>
<td>1389</td>
</tr>
<tr>
<td>ALTER LARGE OBJECT</td>
<td>1390</td>
</tr>
<tr>
<td>ALTER MATERIALIZED VIEW</td>
<td>1391</td>
</tr>
<tr>
<td>ALTER OPERATOR</td>
<td>1393</td>
</tr>
<tr>
<td>ALTER OPERATOR CLASS</td>
<td>1395</td>
</tr>
<tr>
<td>ALTER OPERATOR FAMILY</td>
<td>1396</td>
</tr>
<tr>
<td>ALTER POLICY</td>
<td>1400</td>
</tr>
<tr>
<td>ALTER PROCEDURE</td>
<td>1402</td>
</tr>
<tr>
<td>ALTER PUBLICATION</td>
<td>1405</td>
</tr>
<tr>
<td>ALTER ROLE</td>
<td>1407</td>
</tr>
<tr>
<td>ALTER ROUTINE</td>
<td>1411</td>
</tr>
<tr>
<td>ALTER RULE</td>
<td>1413</td>
</tr>
<tr>
<td>ALTER SCHEMA</td>
<td>1414</td>
</tr>
<tr>
<td>ALTER SEQUENCE</td>
<td>1415</td>
</tr>
<tr>
<td>ALTER SERVER</td>
<td>1418</td>
</tr>
<tr>
<td>ALTER STATISTICS</td>
<td>1420</td>
</tr>
<tr>
<td>ALTER SUBSCRIPTION</td>
<td>1421</td>
</tr>
<tr>
<td>ALTER SYSTEM</td>
<td>1423</td>
</tr>
<tr>
<td>ALTER TABLE</td>
<td>1425</td>
</tr>
<tr>
<td>ALTER TABLESPACE</td>
<td>1441</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH CONFIGURATION</td>
<td>1443</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH DICTIONARY</td>
<td>1445</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH PARSER</td>
<td>1447</td>
</tr>
<tr>
<td>ALTER TEXT SEARCH TEMPLATE</td>
<td>1448</td>
</tr>
<tr>
<td>ALTER TRIGGER</td>
<td>1449</td>
</tr>
<tr>
<td>ALTER TYPE</td>
<td>1451</td>
</tr>
<tr>
<td>ALTER USER</td>
<td>1455</td>
</tr>
<tr>
<td>ALTER USER MAPPING</td>
<td>1456</td>
</tr>
<tr>
<td>ALTER VIEW</td>
<td>1457</td>
</tr>
<tr>
<td>ANALYZE</td>
<td>1459</td>
</tr>
<tr>
<td>BEGIN</td>
<td>1462</td>
</tr>
<tr>
<td>CALL</td>
<td>1464</td>
</tr>
<tr>
<td>CHECKPOINT</td>
<td>1465</td>
</tr>
<tr>
<td>CLOSE</td>
<td>1466</td>
</tr>
<tr>
<td>SQL Commands</td>
<td>Page</td>
</tr>
<tr>
<td>--------------------------------------------------</td>
<td>------</td>
</tr>
<tr>
<td>CLUSTER</td>
<td>1467</td>
</tr>
<tr>
<td>COMMENT</td>
<td>1469</td>
</tr>
<tr>
<td>COMMIT</td>
<td>1474</td>
</tr>
<tr>
<td>COMMIT PREPARED</td>
<td>1475</td>
</tr>
<tr>
<td>COPY</td>
<td>1476</td>
</tr>
<tr>
<td>CREATE ACCESS METHOD</td>
<td>1486</td>
</tr>
<tr>
<td>CREATE AGGREGATE</td>
<td>1487</td>
</tr>
<tr>
<td>CREATE CAST</td>
<td>1495</td>
</tr>
<tr>
<td>CREATE COLLATION</td>
<td>1499</td>
</tr>
<tr>
<td>CREATE CONVERSION</td>
<td>1501</td>
</tr>
<tr>
<td>CREATE DATABASE</td>
<td>1503</td>
</tr>
<tr>
<td>CREATE DOMAIN</td>
<td>1506</td>
</tr>
<tr>
<td>CREATE EVENT TRIGGER</td>
<td>1509</td>
</tr>
<tr>
<td>CREATE EXTENSION</td>
<td>1511</td>
</tr>
<tr>
<td>CREATE FOREIGN DATA WRAPPER</td>
<td>1513</td>
</tr>
<tr>
<td>CREATE FOREIGN TABLE</td>
<td>1515</td>
</tr>
<tr>
<td>CREATE FUNCTION</td>
<td>1519</td>
</tr>
<tr>
<td>CREATE GROUP</td>
<td>1527</td>
</tr>
<tr>
<td>CREATE INDEX</td>
<td>1528</td>
</tr>
<tr>
<td>CREATE LANGUAGE</td>
<td>1536</td>
</tr>
<tr>
<td>CREATE MATERIALIZED VIEW</td>
<td>1539</td>
</tr>
<tr>
<td>CREATE OPERATOR</td>
<td>1541</td>
</tr>
<tr>
<td>CREATE OPERATOR CLASS</td>
<td>1544</td>
</tr>
<tr>
<td>CREATE OPERATOR FAMILY</td>
<td>1547</td>
</tr>
<tr>
<td>CREATE POLICY</td>
<td>1548</td>
</tr>
<tr>
<td>CREATE PROCEDURE</td>
<td>1554</td>
</tr>
<tr>
<td>CREATE PUBLICATION</td>
<td>1557</td>
</tr>
<tr>
<td>CREATE ROLE</td>
<td>1559</td>
</tr>
<tr>
<td>CREATE RULE</td>
<td>1564</td>
</tr>
<tr>
<td>CREATE SCHEMA</td>
<td>1567</td>
</tr>
<tr>
<td>CREATE SEQUENCE</td>
<td>1570</td>
</tr>
<tr>
<td>CREATE SERVER</td>
<td>1574</td>
</tr>
<tr>
<td>CREATE STATISTICS</td>
<td>1576</td>
</tr>
<tr>
<td>CREATE SUBSCRIPTION</td>
<td>1578</td>
</tr>
<tr>
<td>CREATE TABLE</td>
<td>1581</td>
</tr>
<tr>
<td>CREATE TABLE AS</td>
<td>1601</td>
</tr>
<tr>
<td>CREATE TABLESPACE</td>
<td>1604</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH CONFIGURATION</td>
<td>1606</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH DICTIONARY</td>
<td>1607</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH PARSEER</td>
<td>1609</td>
</tr>
<tr>
<td>CREATE TEXT SEARCH TEMPLATE</td>
<td>1611</td>
</tr>
<tr>
<td>CREATE TRANSFORM</td>
<td>1612</td>
</tr>
<tr>
<td>CREATE TRIGGER</td>
<td>1614</td>
</tr>
<tr>
<td>CREATE TYPE</td>
<td>1621</td>
</tr>
<tr>
<td>CREATE USER</td>
<td>1630</td>
</tr>
<tr>
<td>CREATE USER MAPPING</td>
<td>1631</td>
</tr>
<tr>
<td>CREATE VIEW</td>
<td>1633</td>
</tr>
<tr>
<td>DEALLOCATE</td>
<td>1638</td>
</tr>
<tr>
<td>DECLARE</td>
<td>1639</td>
</tr>
<tr>
<td>DELETE</td>
<td>1642</td>
</tr>
<tr>
<td>DISCARD</td>
<td>1645</td>
</tr>
<tr>
<td>DO</td>
<td>1646</td>
</tr>
<tr>
<td>DROP ACCESS METHOD</td>
<td>1648</td>
</tr>
<tr>
<td>DROP AGGREGATE</td>
<td>1649</td>
</tr>
<tr>
<td>DROP CAST</td>
<td>1651</td>
</tr>
<tr>
<td>DROP COLLATION</td>
<td>1652</td>
</tr>
<tr>
<td>DROP CONVERSION</td>
<td>1653</td>
</tr>
<tr>
<td>DROP DATABASE</td>
<td>1654</td>
</tr>
<tr>
<td>SQL Command</td>
<td>Page</td>
</tr>
<tr>
<td>-----------------------------------</td>
<td>------</td>
</tr>
<tr>
<td>DROP DOMAIN</td>
<td>1655</td>
</tr>
<tr>
<td>DROP EVENT TRIGGER</td>
<td>1656</td>
</tr>
<tr>
<td>DROP EXTENSION</td>
<td>1657</td>
</tr>
<tr>
<td>DROP FOREIGN DATA WRAPPER</td>
<td>1658</td>
</tr>
<tr>
<td>DROP FOREIGN TABLE</td>
<td>1659</td>
</tr>
<tr>
<td>DROP FUNCTION</td>
<td>1660</td>
</tr>
<tr>
<td>DROP GROUP</td>
<td>1662</td>
</tr>
<tr>
<td>DROP INDEX</td>
<td>1663</td>
</tr>
<tr>
<td>DROP LANGUAGE</td>
<td>1665</td>
</tr>
<tr>
<td>DROP MATERIALIZED VIEW</td>
<td>1666</td>
</tr>
<tr>
<td>DROP OPERATOR</td>
<td>1667</td>
</tr>
<tr>
<td>DROP OPERATOR CLASS</td>
<td>1669</td>
</tr>
<tr>
<td>DROP OPERATOR FAMILY</td>
<td>1671</td>
</tr>
<tr>
<td>DROP OWNED</td>
<td>1673</td>
</tr>
<tr>
<td>DROP POLICY</td>
<td>1674</td>
</tr>
<tr>
<td>DROP PROCEDURE</td>
<td>1675</td>
</tr>
<tr>
<td>DROP PUBLICATION</td>
<td>1677</td>
</tr>
<tr>
<td>DROP ROLE</td>
<td>1678</td>
</tr>
<tr>
<td>DROP ROUTINE</td>
<td>1679</td>
</tr>
<tr>
<td>DROP RULE</td>
<td>1680</td>
</tr>
<tr>
<td>DROP SCHEMA</td>
<td>1681</td>
</tr>
<tr>
<td>DROP SEQUENCE</td>
<td>1682</td>
</tr>
<tr>
<td>DROP SERVER</td>
<td>1683</td>
</tr>
<tr>
<td>DROP STATISTICS</td>
<td>1684</td>
</tr>
<tr>
<td>DROP SUBSCRIPTION</td>
<td>1685</td>
</tr>
<tr>
<td>DROP TABLE</td>
<td>1687</td>
</tr>
<tr>
<td>DROP TABLESPACE</td>
<td>1688</td>
</tr>
<tr>
<td>DROP TEXT SEARCH CONFIGURATION</td>
<td>1689</td>
</tr>
<tr>
<td>DROP TEXT SEARCH DICTIONARY</td>
<td>1690</td>
</tr>
<tr>
<td>DROP TEXT SEARCH Parser</td>
<td>1691</td>
</tr>
<tr>
<td>DROP TEXT SEARCH TEMPLATE</td>
<td>1692</td>
</tr>
<tr>
<td>DROP TRANSFORM</td>
<td>1693</td>
</tr>
<tr>
<td>DROP TRIGGER</td>
<td>1694</td>
</tr>
<tr>
<td>DROP TYPE</td>
<td>1695</td>
</tr>
<tr>
<td>DROP USER</td>
<td>1696</td>
</tr>
<tr>
<td>DROP USER MAPPING</td>
<td>1697</td>
</tr>
<tr>
<td>DROP VIEW</td>
<td>1698</td>
</tr>
<tr>
<td>END</td>
<td>1699</td>
</tr>
<tr>
<td>EXECUTE</td>
<td>1700</td>
</tr>
<tr>
<td>EXPLAIN</td>
<td>1701</td>
</tr>
<tr>
<td>FETCH</td>
<td>1706</td>
</tr>
<tr>
<td>GRANT</td>
<td>1710</td>
</tr>
<tr>
<td>IMPORT FOREIGN SCHEMA</td>
<td>1717</td>
</tr>
<tr>
<td>INSERT</td>
<td>1719</td>
</tr>
<tr>
<td>LISTEN</td>
<td>1726</td>
</tr>
<tr>
<td>LOAD</td>
<td>1728</td>
</tr>
<tr>
<td>LOCK</td>
<td>1729</td>
</tr>
<tr>
<td>MOVE</td>
<td>1732</td>
</tr>
<tr>
<td>NOTIFY</td>
<td>1734</td>
</tr>
<tr>
<td>PREPARE</td>
<td>1737</td>
</tr>
<tr>
<td>PREPARE TRANSACTION</td>
<td>1740</td>
</tr>
<tr>
<td>REASSIGN OWNED</td>
<td>1742</td>
</tr>
<tr>
<td>REFRESH MATERIALIZED VIEW</td>
<td>1743</td>
</tr>
<tr>
<td>REINDEX</td>
<td>1745</td>
</tr>
<tr>
<td>RELEASE SAVEPOINT</td>
<td>1748</td>
</tr>
<tr>
<td>RESET</td>
<td>1750</td>
</tr>
<tr>
<td>REVOKE</td>
<td>1751</td>
</tr>
<tr>
<td>ROLLBACK</td>
<td>1755</td>
</tr>
</tbody>
</table>
SQL Commands

ROLLBACK PREPARED .................................................................................. 1756
ROLLBACK TO SAVEPOINT ........................................................................ 1757
SAVEPOINT ................................................................................................. 1759
SECURITY LABEL ....................................................................................... 1761
SELECT ........................................................................................................ 1764
SELECT INTO ............................................................................................... 1785
SET ............................................................................................................... 1787
SET CONSTRAINTS ....................................................................................... 1790
SET ROLE ..................................................................................................... 1791
SET SESSION AUTHORIZATION .................................................................. 1793
SET TRANSACTION ...................................................................................... 1795
SHOW ........................................................................................................... 1798
START TRANSACTION ................................................................................ 1800
TRUNCATE ................................................................................................. 1801
UNLISTEN .................................................................................................... 1803
UPDATE ........................................................................................................ 1805
VACUUM ....................................................................................................... 1810
VALUES ....................................................................................................... 1813
ABORT

ABORT — abort the current transaction

Synopsis

ABORT [ WORK | TRANSACTION ]

Description

ABORT rolls back the current transaction and causes all the updates made by the transaction to be discarded. This command is identical in behavior to the standard SQL command ROLLBACK, and is present only for historical reasons.

Parameters

WORK
TRANSACTION

Optional key words. They have no effect.

Notes

Use COMMIT to successfully terminate a transaction.

Issuing ABORT outside of a transaction block emits a warning and otherwise has no effect.

Examples

To abort all changes:

ABORT;

Compatibility

This command is a PostgreSQL extension present for historical reasons. ROLLBACK is the equivalent standard SQL command.

See Also

BEGIN, COMMIT, ROLLBACK
ALTER AGGREGATE

ALTER AGGREGATE — change the definition of an aggregate function

Synopsis

ALTER AGGREGATE name ( aggregate_signature ) RENAME TO new_name
ALTER AGGREGATE name ( aggregate_signature )
           OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER AGGREGATE name ( aggregate_signature ) SET SCHEMA new_schema

where aggregate_signature is:

  * | [ argmode ] [ argname ] argtype [ , ... ] |
  [ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ]
  [ argname ] argtype [ , ... ]

Description

ALTER AGGREGATE changes the definition of an aggregate function.

You must own the aggregate function to use ALTER AGGREGATE. To change the schema of an aggregate function, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the aggregate function's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the aggregate function. However, a superuser can alter ownership of any aggregate function anyway.)

Parameters

name

The name (optionally schema-qualified) of an existing aggregate function.

argmode

The mode of an argument: IN or VARIADIC. If omitted, the default is IN.

argname

The name of an argument. Note that ALTER AGGREGATE does not actually pay any attention to argument names, since only the argument data types are needed to determine the aggregate function's identity.

argtype

An input data type on which the aggregate function operates. To reference a zero-argument aggregate function, write * in place of the list of argument specifications. To reference an ordered-set aggregate function, write ORDER BY between the direct and aggregated argument specifications.

new_name

The new name of the aggregate function.
ALTER AGGREGATE

new_owner

The new owner of the aggregate function.

new_schema

The new schema for the aggregate function.

Notes

The recommended syntax for referencing an ordered-set aggregate is to write ORDER BY between the direct and aggregated argument specifications, in the same style as in CREATE AGGREGATE. However, it will also work to omit ORDER BY and just run the direct and aggregated argument specifications into a single list. In this abbreviated form, if VARIADIC "any" was used in both the direct and aggregated argument lists, write VARIADIC "any" only once.

Examples

To rename the aggregate function myavg for type integer to my_average:

ALTER AGGREGATE myavg(integer) RENAME TO my_average;

To change the owner of the aggregate function myavg for type integer to joe:

ALTER AGGREGATE myavg(integer) OWNER TO joe;

To move the ordered-set aggregate mypercentile with direct argument of type float8 and aggregated argument of type integer into schema myschema:

ALTER AGGREGATE mypercentile(float8 ORDER BY integer) SET SCHEMA myschema;

This will work too:

ALTER AGGREGATE mypercentile(float8, integer) SET SCHEMA myschema;

Compatibility

There is no ALTER AGGREGATE statement in the SQL standard.

See Also

CREATE AGGREGATE, DROP AGGREGATE
ALTER COLLATION

ALTER COLLATION — change the definition of a collation

Synopsis

ALTER COLLATION name REFRESH VERSION

ALTER COLLATION name RENAME TO new_name

ALTER COLLATION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER COLLATION name SET SCHEMA new_schema

Description

ALTER COLLATION changes the definition of a collation.

You must own the collation to use ALTER COLLATION. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the collation's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the collation. However, a superuser can alter ownership of any collation anyway.)

Parameters

name

The name (optionally schema-qualified) of an existing collation.

new_name

The new name of the collation.

new_owner

The new owner of the collation.

new_schema

The new schema for the collation.

REFRESH VERSION

Update the collation's version. See Notes below.

Notes

When using collations provided by the ICU library, the ICU-specific version of the collator is recorded in the system catalog when the collation object is created. When the collation is used, the current version is checked against the recorded version, and a warning is issued when there is a mismatch, for example:

WARNING:  collation "xx-x-icu" has version mismatch
DETAIL:  The collation in the database was created using version 1.2.3.4, but the operating system provides version 2.3.4.5.
HINT: Rebuild all objects affected by this collation and run
ALTER COLLATION pg_catalog."xx-x-icu" REFRESH VERSION, or build
PostgreSQL with the right library version.

A change in collation definitions can lead to corrupt indexes and other problems because the database
system relies on stored objects having a certain sort order. Generally, this should be avoided, but it can
happen in legitimate circumstances, such as when using pg_upgrade to upgrade to server binaries
linked with a newer version of ICU. When this happens, all objects depending on the collation should
be rebuilt, for example, using REINDEX. When that is done, the collation version can be refreshed
using the command ALTER COLLATION ... REFRESH VERSION. This will update the system
catalog to record the current collator version and will make the warning go away. Note that this does
not actually check whether all affected objects have been rebuilt correctly.

The following query can be used to identify all collations in the current database that need to be
refreshed and the objects that depend on them:

```
SELECT pg_describe_object(refclassid, refobjid, refobjsubid) AS "Collation",
       pg_describe_object(classid, objid, objsubid) AS "Object"
FROM pg_depend d JOIN pg_collation c
     ON refclassid = 'pg_collation'::regclass AND refobjid =
       c.oid
WHERE c.collversion <> pg_collation_actual_version(c.oid)
ORDER BY 1, 2;
```

**Examples**

To rename the collation de_DE to german:

```
ALTER COLLATION "de_DE" RENAME TO german;
```

To change the owner of the collation en_US to joe:

```
ALTER COLLATION "en_US" OWNER TO joe;
```

**Compatibility**

There is no ALTER COLLATION statement in the SQL standard.

**See Also**

CREATE COLLATION, DROP COLLATION
ALTER CONVERSION

ALTER CONVERSION — change the definition of a conversion

Synopsis

ALTER CONVERSION name RENAME TO new_name
ALTER CONVERSION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER CONVERSION name SET SCHEMA new_schema

Description

ALTER CONVERSION changes the definition of a conversion.

You must own the conversion to use ALTER CONVERSION. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the conversion’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the conversion. However, a superuser can alter ownership of any conversion anyway.)

Parameters

name

The name (optionally schema-qualified) of an existing conversion.

new_name

The new name of the conversion.

new_owner

The new owner of the conversion.

new_schema

The new schema for the conversion.

Examples

To rename the conversion iso_8859_1_to_utf8 to latin1_to_unicode:

ALTER CONVERSION iso_8859_1_to_utf8 RENAME TO latin1_to_unicode;

To change the owner of the conversion iso_8859_1_to_utf8 to joe:

ALTER CONVERSION iso_8859_1_to_utf8 OWNER TO joe;

Compatibility

There is no ALTER CONVERSION statement in the SQL standard.
ALTER CONVERSION

See Also
CREATE CONVERSION, DROP CONVERSION
ALTER DATABASE

ALTER DATABASE — change a database

Synopsis

ALTER DATABASE name [ [ WITH ] option [ ... ] ]

where option can be:

   ALLOW_CONNECTIONS allowconn
   CONNECTION LIMIT connlimit
   IS_TEMPLATE istemplate

ALTER DATABASE name RENAME TO new_name

ALTER DATABASE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER DATABASE name SET TABLESPACE new_tablespace

ALTER DATABASE name SET configuration_parameter { TO | = } { value | DEFAULT }

ALTER DATABASE name SET configuration_parameter FROM CURRENT

ALTER DATABASE name RESET configuration_parameter

ALTER DATABASE name RESET ALL

Description

ALTER DATABASE changes the attributes of a database.

The first form changes certain per-database settings. (See below for details.) Only the database owner or a superuser can change these settings.

The second form changes the name of the database. Only the database owner or a superuser can rename a database; non-superuser owners must also have the CREATEDB privilege. The current database cannot be renamed. (Connect to a different database if you need to do that.)

The third form changes the owner of the database. To alter the owner, you must own the database and also be a direct or indirect member of the new owning role, and you must have the CREATEDB privilege. (Note that superusers have all these privileges automatically.)

The fourth form changes the default tablespace of the database. Only the database owner or a superuser can do this; you must also have create privilege for the new tablespace. This command physically moves any tables or indexes in the database's old default tablespace to the new tablespace. The new default tablespace must be empty for this database, and no one can be connected to the database. Tables and indexes in non-default tablespaces are unaffected.

The remaining forms change the session default for a run-time configuration variable for a PostgreSQL database. Whenever a new session is subsequently started in that database, the specified value becomes the session default value. The database-specific default overrides whatever setting is present in postgresql.conf or has been received from the postgres command line. Only the database owner or a superuser can change the session defaults for a database. Certain variables cannot be set this way, or can only be set by a superuser.
Parameters

name

The name of the database whose attributes are to be altered.

allowconn

If false then no one can connect to this database.

connlimit

How many concurrent connections can be made to this database. -1 means no limit.

istemplate

If true, then this database can be cloned by any user with CREATEDB privileges; if false, then only superusers or the owner of the database can clone it.

new_name

The new name of the database.

new_owner

The new owner of the database.

new_tablespace

The new default tablespace of the database.

This form of the command cannot be executed inside a transaction block.

configuration_parameter value

Set this database's session default for the specified configuration parameter to the given value. If value is DEFAULT or, equivalently, RESET is used, the database-specific setting is removed, so the system-wide default setting will be inherited in new sessions. Use RESET ALL to clear all database-specific settings. SET FROM CURRENT saves the session's current value of the parameter as the database-specific value.

See SET and Chapter 19 for more information about allowed parameter names and values.

Notes

It is also possible to tie a session default to a specific role rather than to a database; see ALTER ROLE. Role-specific settings override database-specific ones if there is a conflict.

Examples

To disable index scans by default in the database test:

```
ALTER DATABASE test SET enable_indexscan TO off;
```

Compatibility

The ALTER DATABASE statement is a PostgreSQL extension.
ALTER DATABASE

See Also

CREATE DATABASE, DROP DATABASE, SET, CREATE TABLESPACE
ALTER DEFAULT PRIVILEGES

ALTER DEFAULT PRIVILEGES — define default access privileges

Synopsis

ALTER DEFAULT PRIVILEGES
[ FOR { ROLE | USER } target_role [, ...] ]
[ IN SCHEMA schema_name [, ...] ]
abbreviated_grant_or_revoke

where abbreviated_grant_or_revoke is one of:

GRANT { ( SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES |
TRIGGER )
[, ...] | ALL [ PRIVILEGES ] }
ON TABLES
TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { ( USAGE | SELECT | UPDATE )
[, ...] | ALL [ PRIVILEGES ] }
ON SEQUENCES
TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { EXECUTE | ALL [ PRIVILEGES ] }
ON { FUNCTIONS | ROUTINES }
TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON TYPES
TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | CREATE | ALL [ PRIVILEGES ] }
ON SCHEMAS
TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

REVOKE [ GRANT OPTION FOR ]
( { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES |
TRIGGER )
[, ...] | ALL [ PRIVILEGES ] }
ON TABLES
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
( { USAGE | SELECT | UPDATE }
[, ...] | ALL [ PRIVILEGES ] )
ON SEQUENCES
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
ALTER DEFAULT PRIVILEGES

REVOKE [ GRANT OPTION FOR ]
{ EXECUTE | ALL [ PRIVILEGES ] }
ON { FUNCTIONS | ROUTINES }
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
{ USAGE | ALL [ PRIVILEGES ] }
ON TYPES
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
{ USAGE | CREATE | ALL [ PRIVILEGES ] }
ON SCHEMAS
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

Description

ALTER DEFAULT PRIVILEGES allows you to set the privileges that will be applied to objects created in the future. (It does not affect privileges assigned to already-existing objects.) Currently, only the privileges for schemas, tables (including views and foreign tables), sequences, functions, and types (including domains) can be altered. For this command, functions include aggregates and procedures. The words FUNCTIONS and ROUTINES are equivalent in this command. (ROUTINES is preferred going forward as the standard term for functions and procedures taken together. In earlier PostgreSQL releases, only the word FUNCTIONS was allowed. It is not possible to set default privileges for functions and procedures separately.)

You can change default privileges only for objects that will be created by yourself or by roles that you are a member of. The privileges can be set globally (i.e., for all objects created in the current database), or just for objects created in specified schemas. Default privileges that are specified per-schema are added to whatever the global default privileges are for the particular object type.

As explained under GRANT, the default privileges for any object type normally grant all grantable permissions to the object owner, and may grant some privileges to PUBLIC as well. However, this behavior can be changed by altering the global default privileges with ALTER DEFAULT PRIVILEGES.

Parameters

target_role

The name of an existing role of which the current role is a member. If FOR ROLE is omitted, the current role is assumed.

schema_name

The name of an existing schema. If specified, the default privileges are altered for objects later created in that schema. If IN SCHEMA is omitted, the global default privileges are altered. IN SCHEMA is not allowed when using ON SCHEMAS as schemas can’t be nested.

role_name

The name of an existing role to grant or revoke privileges for. This parameter, and all the other parameters in abbreviated_grant_or_revoke, act as described under GRANT or REVOKE, except that one is setting permissions for a whole class of objects rather than specific named objects.
Notes

Use psql's \ddp command to obtain information about existing assignments of default privileges. The meaning of the privilege values is the same as explained for \dp under GRANT.

If you wish to drop a role for which the default privileges have been altered, it is necessary to reverse the changes in its default privileges or use DROP OWNED BY to get rid of the default privileges entry for the role.

Examples

Grant SELECT privilege to everyone for all tables (and views) you subsequently create in schema myschema, and allow role webuser to INSERT into them too:

```
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema GRANT SELECT ON TABLES TO PUBLIC;
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema GRANT INSERT ON TABLES TO webuser;
```

Undo the above, so that subsequently-created tables won't have any more permissions than normal:

```
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema REVOKE SELECT ON TABLES FROM PUBLIC;
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema REVOKE INSERT ON TABLES FROM webuser;
```

Remove the public EXECUTE permission that is normally granted on functions, for all functions subsequently created by role admin:

```
ALTER DEFAULT PRIVILEGES FOR ROLE admin REVOKE EXECUTE ON FUNCTIONS FROM PUBLIC;
```

Compatibility

There is no ALTER DEFAULT PRIVILEGES statement in the SQL standard.

See Also

GRANT, REVOKE
ALTER DOMAIN

ALTER DOMAIN — change the definition of a domain

Synopsis

```
ALTER DOMAIN name 
   { SET DEFAULT expression | DROP DEFAULT }
ALTER DOMAIN name 
   { SET | DROP } NOT NULL
ALTER DOMAIN name
   ADD domain_constraint [ NOT VALID ]
ALTER DOMAIN name
   DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT | CASCADE ]
ALTER DOMAIN name
   RENAME CONSTRAINT constraint_name TO new_constraint_name
ALTER DOMAIN name
   VALIDATE CONSTRAINT constraint_name
ALTER DOMAIN name
   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER DOMAIN name
   RENAME TO new_name
ALTER DOMAIN name
   SET SCHEMA new_schema
```

Description

ALTER DOMAIN changes the definition of an existing domain. There are several sub-forms:

SET/DROP DEFAULT

These forms set or remove the default value for a domain. Note that defaults only apply to subsequent INSERT commands; they do not affect rows already in a table using the domain.

SET/DROP NOT NULL

These forms change whether a domain is marked to allow NULL values or to reject NULL values. You can only SET NOT NULL when the columns using the domain contain no null values.

ADD domain_constraint [ NOT VALID ]

This form adds a new constraint to a domain using the same syntax as CREATE DOMAIN. When a new constraint is added to a domain, all columns using that domain will be checked against the newly added constraint. These checks can be suppressed by adding the new constraint using the NOT VALID option; the constraint can later be made valid using ALTER DOMAIN ... VALIDATE CONSTRAINT. Newly inserted or updated rows are always checked against all constraints, even those marked NOT VALID. NOT VALID is only accepted for CHECK constraints.

DROP CONSTRAINT [ IF EXISTS ]

This form drops constraints on a domain. If IF EXISTS is specified and the constraint does not exist, no error is thrown. In this case a notice is issued instead.

RENAME CONSTRAINT

This form changes the name of a constraint on a domain.
VALIDATE CONSTRAINT

This form validates a constraint previously added as NOT VALID, that is, verify that all data in columns using the domain satisfy the specified constraint.

OWNER

This form changes the owner of the domain to the specified user.

RENAME

This form changes the name of the domain.

SET SCHEMA

This form changes the schema of the domain. Any constraints associated with the domain are moved into the new schema as well.

You must own the domain to use ALTER DOMAIN. To change the schema of a domain, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the domain's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the domain. However, a superuser can alter ownership of any domain anyway.)

Parameters

name

The name (possibly schema-qualified) of an existing domain to alter.

domain_constraint

New domain constraint for the domain.

constraint_name

Name of an existing constraint to drop or rename.

NOT VALID

Do not verify existing column data for constraint validity.

CASCADE

Automatically drop objects that depend on the constraint, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the constraint if there are any dependent objects. This is the default behavior.

new_name

The new name for the domain.

new_constraint_name

The new name for the constraint.

new_owner

The user name of the new owner of the domain.
The new schema for the domain.

Notes

Currently, ALTER DOMAIN ADD CONSTRAINT, ALTER DOMAIN VALIDATE CONSTRAINT, and ALTER DOMAIN SET NOT NULL will fail if the named domain or any derived domain is used within a container-type column (a composite, array, or range column) in any table in the database. They should eventually be improved to be able to verify the new constraint for such nested values.

Examples

To add a NOT NULL constraint to a domain:

```
ALTER DOMAIN zipcode SET NOT NULL;
```

To remove a NOT NULL constraint from a domain:

```
ALTER DOMAIN zipcode DROP NOT NULL;
```

To add a check constraint to a domain:

```
ALTER DOMAIN zipcode ADD CONSTRAINT zipchk CHECK (char_length(VALUE) = 5);
```

To remove a check constraint from a domain:

```
ALTER DOMAIN zipcode DROP CONSTRAINT zipchk;
```

To rename a check constraint on a domain:

```
ALTER DOMAIN zipcode RENAME CONSTRAINT zipchk TO zip_check;
```

To move the domain into a different schema:

```
ALTER DOMAIN zipcode SET SCHEMA customers;
```

Compatibility

ALTER DOMAIN conforms to the SQL standard, except for the OWNER, RENAME, SET SCHEMA, and VALIDATE CONSTRAINT variants, which are PostgreSQL extensions. The NOT VALID clause of the ADD CONSTRAINT variant is also a PostgreSQL extension.

See Also

CREATE DOMAIN, DROP DOMAIN
ALTER EVENT TRIGGER

ALTER EVENT TRIGGER — change the definition of an event trigger

Synopsis

ALTER EVENT TRIGGER name DISABLE
ALTER EVENT TRIGGER name ENABLE [ REPLICA | ALWAYS ]
ALTER EVENT TRIGGER name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER EVENT TRIGGER name RENAME TO new_name

Description

ALTER EVENT TRIGGER changes properties of an existing event trigger.

You must be superuser to alter an event trigger.

Parameters

name

The name of an existing trigger to alter.

new_owner

The user name of the new owner of the event trigger.

new_name

The new name of the event trigger.

DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER

These forms configure the firing of event triggers. A disabled trigger is still known to the system, but is not executed when its triggering event occurs. See also session_replication_role.

Compatibility

There is no ALTER EVENT TRIGGER statement in the SQL standard.

See Also

CREATE EVENT TRIGGER, DROP EVENT TRIGGER
ALTER EXTENSION

ALTER EXTENSION — change the definition of an extension

Synopsis

ALTER EXTENSION name UPDATE [ TO new_version ]
ALTER EXTENSION name SET SCHEMA new_schema
ALTER EXTENSION name ADD member_object
ALTER EXTENSION name DROP member_object

where member_object is:

ACCESS METHOD object_name |
AGGREGATE aggregate_name ( aggregate_signature ) |
CAST (source_type AS target_type) |
COLLATION object_name |
CONVERSION object_name |
DOMAIN object_name |
EVENT TRIGGER object_name |
FOREIGN DATA WRAPPER object_name |
FOREIGN TABLE object_name |
FUNCTION function_name [ ( ( [ argmode ] [ argname ] argtype [, ... ] ) ) ] |
MATERIALIZED VIEW object_name |
OPERATOR operator_name (left_type, right_type) |
OPERATOR CLASS object_name USING index_method |
OPERATOR FAMILY object_name USING index_method |
[ PROCEDURAL ] LANGUAGE object_name |
PROCEDURE procedure_name [ ( ( [ argmode ] [ argname ] argtype [, ... ] ) ) ] |
ROUTINE routine_name [ ( ( [ argmode ] [ argname ] argtype [, ... ] ) ) ] |
SCHEMA object_name |
SEQUENCE object_name |
SERVER object_name |
TABLE object_name |
TEXT SEARCH CONFIGURATION object_name |
TEXT SEARCH DICTIONARY object_name |
TEXT SEARCH PARSER object_name |
TEXT SEARCH TEMPLATE object_name |
TRANSFORM FOR type_name LANGUAGE lang_name |
TYPE object_name |
VIEW object_name |

and aggregate_signature is:

* |
[ argmode ] [ argname ] argtype [ , ... ] |
[ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ]
[ argname ] argtype [ , ... ]
ALTER EXTENSION

Description

ALTER EXTENSION changes the definition of an installed extension. There are several subforms:

UPDATE

This form updates the extension to a newer version. The extension must supply a suitable update script (or series of scripts) that can modify the currently-installed version into the requested version.

SET SCHEMA

This form moves the extension’s objects into another schema. The extension has to be relocatable for this command to succeed.

ADD member_object

This form adds an existing object to the extension. This is mainly useful in extension update scripts. The object will subsequently be treated as a member of the extension; notably, it can only be dropped by dropping the extension.

DROP member_object

This form removes a member object from the extension. This is mainly useful in extension update scripts. The object is not dropped, only disassociated from the extension.

See Section 38.16 for more information about these operations.

You must own the extension to use ALTER EXTENSION. The ADD/DROP forms require ownership of the added/dropped object as well.

Parameters

name

The name of an installed extension.

new_version

The desired new version of the extension. This can be written as either an identifier or a string literal. If not specified, ALTER EXTENSION UPDATE attempts to update to whatever is shown as the default version in the extension's control file.

new_schema

The new schema for the extension.

object_name
aggregate_name
function_name
operator_name
procedure_name
routine_name

The name of an object to be added to or removed from the extension. Names of tables, aggregates, domains, foreign tables, functions, operators, operator classes, operator families, procedures, routines, sequences, text search objects, types, and views can be schema-qualified.

source_type

The name of the source data type of the cast.
target_type

The name of the target data type of the cast.

argmode

The mode of a function, procedure, or aggregate argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that ALTER EXTENSION does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname

The name of a function, procedure, or aggregate argument. Note that ALTER EXTENSION does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

argtype

The data type of a function, procedure, or aggregate argument.

left_type

right_type

The data type(s) of the operator’s arguments (optionally schema-qualified). Write NONE for the missing argument of a prefix or postfix operator.

Examples

To update the hstore extension to version 2.0:

```
ALTER EXTENSION hstore UPDATE TO '2.0';
```

To change the schema of the hstore extension to utils:

```
ALTER EXTENSION hstore SET SCHEMA utils;
```

To add an existing function to the hstore extension:

```
ALTER EXTENSION hstore ADD FUNCTION populate_record(anyelement, hstore);
```

Compatibility

ALTER EXTENSION is a PostgreSQL extension.
See Also
CREATE EXTENSION, DROP EXTENSION
ALTER FOREIGN DATA WRAPPER

ALTER FOREIGN DATA WRAPPER — change the definition of a foreign-data wrapper

Synopsis

ALTER FOREIGN DATA WRAPPER name
  [ HANDLER handler_function | NO HANDLER ]
  [ VALIDATOR validator_function | NO VALIDATOR ]
  [ OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ]) ]
ALTER FOREIGN DATA WRAPPER name OWNER TO ( new_owner | CURRENT_USER
  | SESSION_USER )
ALTER FOREIGN DATA WRAPPER name RENAME TO new_name

Description

ALTER FOREIGN DATA WRAPPER changes the definition of a foreign-data wrapper. The first form
of the command changes the support functions or the generic options of the foreign-data wrapper (at
least one clause is required). The second form changes the owner of the foreign-data wrapper.

Only superusers can alter foreign-data wrappers. Additionally, only superusers can own foreign-data
wrappers.

Parameters

name

The name of an existing foreign-data wrapper.

HANDLER handler_function

Specifies a new handler function for the foreign-data wrapper.

NO HANDLER

This is used to specify that the foreign-data wrapper should no longer have a handler function.

Note that foreign tables that use a foreign-data wrapper with no handler cannot be accessed.

VALIDATOR validator_function

Specifies a new validator function for the foreign-data wrapper.

Note that it is possible that pre-existing options of the foreign-data wrapper, or of dependent
servers, user mappings, or foreign tables, are invalid according to the new validator. PostgreSQL
does not check for this. It is up to the user to make sure that these options are correct before using
the modified foreign-data wrapper. However, any options specified in this ALTER FOREIGN
DATA WRAPPER command will be checked using the new validator.

NO VALIDATOR

This is used to specify that the foreign-data wrapper should no longer have a validator function.

OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )

Change options for the foreign-data wrapper. ADD, SET, and DROP specify the action to be per-
formed. ADD is assumed if no operation is explicitly specified. Option names must be unique;
names and values are also validated using the foreign data wrapper's validator function, if any.
new_owner

The user name of the new owner of the foreign-data wrapper.

new_name

The new name for the foreign-data wrapper.

Examples

Change a foreign-data wrapper dbi, add option foo, drop bar:

ALTER FOREIGN DATA WRAPPER dbi OPTIONS (ADD foo '1', DROP 'bar');

Change the foreign-data wrapper dbi validator to bob.myvalidator:

ALTER FOREIGN DATA WRAPPER dbi VALIDATOR bob.myvalidator;

Compatibility

ALTER FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED), except that the HANDLER, VALIDATOR, OWNER TO, and RENAME clauses are extensions.

See Also

CREATE FOREIGN DATA WRAPPER, DROP FOREIGN DATA WRAPPER
ALTER FOREIGN TABLE — change the definition of a foreign table

Synopsis

```
ALTER FOREIGN TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
  action [, ... ]
ALTER FOREIGN TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
  RENAME [ COLUMN ] column_name TO new_column_name
ALTER FOREIGN TABLE [ IF EXISTS ] name
  RENAME TO new_name
ALTER FOREIGN TABLE [ IF EXISTS ] name
  SET SCHEMA new_schema
```

where action is one of:

```
ADD [ COLUMN ] column_name data_type [ COLLATE collation ]
  [ column_constraint [ ... ] ]
DROP [ COLUMN ] [ IF EXISTS ] column_name [ RESTRICT ]
CASCADE
ALTER [ COLUMN ] column_name [ SET DATA ] TYPE data_type
  [ COLLATE collation ]
ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER [ COLUMN ] column_name { SET | DROP } NOT NULL
ALTER [ COLUMN ] column_name SET STATISTICS integer
ALTER [ COLUMN ] column_name SET ( attribute_option = value
  [, ... ] )
ALTER [ COLUMN ] column_name RESET ( attribute_option
  [, ... ] )
ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL |
  MAIN }
ALTER [ COLUMN ] column_name OPTIONS ( [ ADD | SET |
  DROP ] option ['value'] [, ... ] )
ADD table_constraint [ NOT VALID ]
VALIDATE CONSTRAINT constraint_name
DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT ]
CASCADE
DISABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE REPLICA TRIGGER trigger_name
ENABLE ALWAYS TRIGGER trigger_name
SET WITH OIDS
SET WITHOUT OIDS
INHERIT parent_table
NO INHERIT parent_table
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )
```

Description

ALTER FOREIGN TABLE changes the definition of an existing foreign table. There are several subforms:
ADD COLUMN

This form adds a new column to the foreign table, using the same syntax as CREATE FOREIGN TABLE. Unlike the case when adding a column to a regular table, nothing happens to the underlying storage: this action simply declares that some new column is now accessible through the foreign table.

DROP COLUMN [ IF EXISTS ]

This form drops a column from a foreign table. You will need to say CASCADE if anything outside the table depends on the column; for example, views. If IF EXISTS is specified and the column does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of a column of a foreign table. Again, this has no effect on any underlying storage: this action simply changes the type that PostgreSQL believes the column to have.

SET/DROP DEFAULT

These forms set or remove the default value for a column. Default values only apply in subsequent INSERT or UPDATE commands; they do not cause rows already in the table to change.

SET/DROP NOT NULL

Mark a column as allowing, or not allowing, null values.

SET STATISTICS

This form sets the per-column statistics-gathering target for subsequent ANALYZE operations. See the similar form of ALTER TABLE for more details.

SET ( attribute_option = value [, ... ] )
RESET ( attribute_option [, ... ] )

This form sets or resets per-attribute options. See the similar form of ALTER TABLE for more details.

SET STORAGE

This form sets the storage mode for a column. See the similar form of ALTER TABLE for more details. Note that the storage mode has no effect unless the table's foreign-data wrapper chooses to pay attention to it.

ADD table_constraint [ NOT VALID ]

This form adds a new constraint to a foreign table, using the same syntax as CREATE FOREIGN TABLE. Currently only CHECK constraints are supported.

Unlike the case when adding a constraint to a regular table, nothing is done to verify the constraint is correct; rather, this action simply declares that some new condition should be assumed to hold for all rows in the foreign table. (See the discussion in CREATE FOREIGN TABLE.) If the constraint is marked NOT VALID, then it isn't assumed to hold, but is only recorded for possible future use.

VALIDATE CONSTRAINT

This form marks as valid a constraint that was previously marked as NOT VALID. No action is taken to verify the constraint, but future queries will assume that it holds.

DROP CONSTRAINT [ IF EXISTS ]

This form drops the specified constraint on a foreign table. If IF EXISTS is specified and the constraint does not exist, no error is thrown. In this case a notice is issued instead.
ALTER FOREIGN TABLE

DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER

These forms configure the firing of trigger(s) belonging to the foreign table. See the similar form of ALTER TABLE for more details.

SET WITH OIDS

This form adds an oid system column to the table (see Section 5.4). It does nothing if the table already has OIDs. Unless the table's foreign-data wrapper supports OIDs, this column will simply read as zeroes.

Note that this is not equivalent to ADD COLUMN oid oid; that would add a normal column that happened to be named oid, not a system column.

SET WITHOUT OIDS

This form removes the oid system column from the table. This is exactly equivalent to DROP COLUMN oid RESTRICT, except that it will not complain if there is already no oid column.

INHERIT parent_table

This form adds the target foreign table as a new child of the specified parent table. See the similar form of ALTER TABLE for more details.

NO INHERIT parent_table

This form removes the target foreign table from the list of children of the specified parent table.

OWNER

This form changes the owner of the foreign table to the specified user.

OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )

Change options for the foreign table or one of its columns. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Duplicate option names are not allowed (although it's OK for a table option and a column option to have the same name). Option names and values are also validated using the foreign data wrapper library.

RENAME

The RENAME forms change the name of a foreign table or the name of an individual column in a foreign table.

SET SCHEMA

This form moves the foreign table into another schema.

All the actions except RENAME and SET SCHEMA can be combined into a list of multiple alterations to apply in parallel. For example, it is possible to add several columns and/or alter the type of several columns in a single command.

If the command is written as ALTER FOREIGN TABLE IF EXISTS ... and the foreign table does not exist, no error is thrown. A notice is issued in this case.

You must own the table to use ALTER FOREIGN TABLE. To change the schema of a foreign table, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the table's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the table. However, a superuser can alter ownership of any table anyway.) To add a column or alter a column type, you must also have USAGE privilege on the data type.
Parameters

name

The name (possibly schema-qualified) of an existing foreign table to alter. If ONLY is specified before the table name, only that table is altered. If ONLY is not specified, the table and all its descendant tables (if any) are altered. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

column_name

Name of a new or existing column.

new_column_name

New name for an existing column.

new_name

New name for the table.

data_type

Data type of the new column, or new data type for an existing column.

table_constraint

New table constraint for the foreign table.

constraint_name

Name of an existing constraint to drop.

CASCADE

Automatically drop objects that depend on the dropped column or constraint (for example, views referencing the column), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the column or constraint if there are any dependent objects. This is the default behavior.

trigger_name

Name of a single trigger to disable or enable.

ALL

Disable or enable all triggers belonging to the foreign table. (This requires superuser privilege if any of the triggers are internally generated triggers. The core system does not add such triggers to foreign tables, but add-on code could do so.)

USER

Disable or enable all triggers belonging to the foreign table except for internally generated triggers.

parent_table

A parent table to associate or de-associate with this foreign table.
ALTER FOREIGN TABLE

new_owner

The user name of the new owner of the table.

new_schema

The name of the schema to which the table will be moved.

Notes

The key word COLUMN is noise and can be omitted.

Consistency with the foreign server is not checked when a column is added or removed with ADD COLUMN or DROP COLUMN, a NOT NULL or CHECK constraint is added, or a column type is changed with SET DATA TYPE. It is the user's responsibility to ensure that the table definition matches the remote side.

Refer to CREATE FOREIGN TABLE for a further description of valid parameters.

Examples

To mark a column as not-null:

ALTER FOREIGN TABLE distributors ALTER COLUMN street SET NOT NULL;

To change options of a foreign table:

ALTER FOREIGN TABLE myschema.distributors OPTIONS (ADD opt1 'value', SET opt2 'value2', DROP opt3 'value3');

Compatibility

The forms ADD, DROP, and SET DATA TYPE conform with the SQL standard. The other forms are PostgreSQL extensions of the SQL standard. Also, the ability to specify more than one manipulation in a single ALTER FOREIGN TABLE command is an extension.

ALTER FOREIGN TABLE DROP COLUMN can be used to drop the only column of a foreign table, leaving a zero-column table. This is an extension of SQL, which disallows zero-column foreign tables.

See Also

CREATE FOREIGN TABLE, DROP FOREIGN TABLE
ALTER FUNCTION

ALTER FUNCTION — change the definition of a function

Synopsis

```
ALTER FUNCTION name [ ( [ argmode ] [ argname ] argtype [, ...] ) ]
    action [ ... ] [ RESTRICT ]
ALTER FUNCTION name [ ( [ argmode ] [ argname ] argtype [, ...] ) ]
    RENAME TO new_name
ALTER FUNCTION name [ ( [ argmode ] [ argname ] argtype [, ...] ) ]
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER FUNCTION name [ ( [ argmode ] [ argname ] argtype [, ...] ) ]
    SET SCHEMA new_schema
ALTER FUNCTION name [ ( [ argmode ] [ argname ] argtype [, ...] ) ]
    DEPENDS ON EXTENSION extension_name
```

where `action` is one of:

```
CALLED ON NULL INPUT | RETURNS NULL ON NULL INPUT | STRICT
IMMUTABLE | STABLE | VOLATILE | [ NOT ] LEAKPROOF
[ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
PARALLEL { UNSAFE | RESTRICTED | SAFE }
COST execution_cost
ROWS result_rows
SET configuration_parameter { TO | = } { value | DEFAULT }
SET configuration_parameter FROM CURRENT
RESET configuration_parameter
RESET ALL
```

Description

ALTER FUNCTION changes the definition of a function.

You must own the function to use ALTER FUNCTION. To change a function's schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the function's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the function. However, a superuser can alter ownership of any function anyway.)

Parameters

`name`

The name (optionally schema-qualified) of an existing function. If no argument list is specified, the name must be unique in its schema.

`argmode`

The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that ALTER FUNCTION does not actually pay any attention to OUT arguments, since only the
input arguments are needed to determine the function’s identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname

The name of an argument. Note that ALTER FUNCTION does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

argtype

The data type(s) of the function’s arguments (optionally schema-qualified), if any.

new_name

The new name of the function.

new_owner

The new owner of the function. Note that if the function is marked SECURITY DEFINER, it will subsequently execute as the new owner.

new_schema

The new schema for the function.

extension_name

The name of the extension that the function is to depend on.

CALLED ON NULL INPUT
RETURNS NULL ON NULL INPUT
STRICT

CALLED ON NULL INPUT changes the function so that it will be invoked when some or all of its arguments are null. RETURNS NULL ON NULL INPUT or STRICT changes the function so that it is not invoked if any of its arguments are null; instead, a null result is assumed automatically. See CREATE FUNCTION for more information.

IMMUTABLE
STABLE
VOLATILE

Change the volatility of the function to the specified setting. See CREATE FUNCTION for details.

[ EXTERNAL ] SECURITY INVOKER
[ EXTERNAL ] SECURITY DEFINER

Change whether the function is a security definer or not. The key word EXTERNAL is ignored for SQL conformance. See CREATE FUNCTION for more information about this capability.

PARALLEL

Change whether the function is deemed safe for parallelism. See CREATE FUNCTION for details.

LEAKPROOF

Change whether the function is considered leakproof or not. See CREATE FUNCTION for more information about this capability.
ALTER FUNCTION

COST execution_cost

Change the estimated execution cost of the function. See CREATE FUNCTION for more information.

ROWS result_rows

Change the estimated number of rows returned by a set-returning function. See CREATE FUNCTION for more information.

configuration_parameter value

Add or change the assignment to be made to a configuration parameter when the function is called. If value is DEFAULT or, equivalently, RESET is used, the function-local setting is removed, so that the function executes with the value present in its environment. Use RESET ALL to clear all function-local settings. SET FROM CURRENT saves the value of the parameter that is current when ALTER FUNCTION is executed as the value to be applied when the function is entered.

See SET and Chapter 19 for more information about allowed parameter names and values.

RESTRICT

Ignored for conformance with the SQL standard.

Examples

To rename the function sqrt for type integer to square_root:

ALTER FUNCTION sqrt(integer) RENAME TO square_root;

To change the owner of the function sqrt for type integer to joe:

ALTER FUNCTION sqrt(integer) OWNER TO joe;

To change the schema of the function sqrt for type integer to maths:

ALTER FUNCTION sqrt(integer) SET SCHEMA maths;

To mark the function sqrt for type integer as being dependent on the extension mathlib:

ALTER FUNCTION sqrt(integer) DEPENDS ON EXTENSION mathlib;

To adjust the search path that is automatically set for a function:

ALTER FUNCTION check_password(text) SET search_path = admin, pg_temp;

To disable automatic setting of search_path for a function:

ALTER FUNCTION check_password(text) RESET search_path;

The function will now execute with whatever search path is used by its caller.
Compatibility

This statement is partially compatible with the ALTER FUNCTION statement in the SQL standard. The standard allows more properties of a function to be modified, but does not provide the ability to rename a function, make a function a security definer, attach configuration parameter values to a function, or change the owner, schema, or volatility of a function. The standard also requires the RESTRICT key word, which is optional in PostgreSQL.

See Also

CREATE FUNCTION, DROP FUNCTION, ALTER PROCEDURE, ALTER ROUTINE
ALTER GROUP

ALTER GROUP — change role name or membership

Synopsis

ALTER GROUP role_specification ADD USER user_name [, ... ]
ALTER GROUP role_specification DROP USER user_name [, ... ]

where role_specification can be:

    role_name
    | CURRENT_USER
    | SESSION_USER

ALTER GROUP group_name RENAME TO new_name

Description

ALTER GROUP changes the attributes of a user group. This is an obsolete command, though still accepted for backwards compatibility, because groups (and users too) have been superseded by the more general concept of roles.

The first two variants add users to a group or remove them from a group. (Any role can play the part of either a “user” or a “group” for this purpose.) These variants are effectively equivalent to granting or revoking membership in the role named as the “group”; so the preferred way to do this is to use GRANT or REVOKE.

The third variant changes the name of the group. This is exactly equivalent to renaming the role with ALTER ROLE.

Parameters

  group_name

    The name of the group (role) to modify.

  user_name

    Users (roles) that are to be added to or removed from the group. The users must already exist; ALTER GROUP does not create or drop users.

  new_name

    The new name of the group.

Examples

Add users to a group:

    ALTER GROUP staff ADD USER karl, john;

Remove a user from a group:
ALTER GROUP workers DROP USER beth;

Compatibility

There is no ALTER GROUP statement in the SQL standard.

See Also

GRANT, REVOKE, ALTER ROLE
ALTER INDEX — change the definition of an index

Synopsis

ALTER INDEX [ IF EXISTS ] name RENAME TO new_name
ALTER INDEX [ IF EXISTS ] name SET TABLESPACE tablespace_name
ALTER INDEX name ATTACH PARTITION index_name
ALTER INDEX name DEPENDS ON EXTENSION extension_name
ALTER INDEX [ IF EXISTS ] name SET ( storage_parameter = value [, ... ] )
ALTER INDEX [ IF EXISTS ] name RESET ( storage_parameter [, ... ] )
ALTER INDEX [ IF EXISTS ] name ALTER [ COLUMN ] column_number
  SET STATISTICS integer
ALTER INDEX ALL IN TABLESPACE name [ OWNED BY role_name [, ... ] ]
  SET TABLESPACE new_tablespace [ NOWAIT ]

Description

ALTER INDEX changes the definition of an existing index. There are several subforms:

RENAME

The RENAME form changes the name of the index. If the index is associated with a table constraint (either UNIQUE, PRIMARY KEY, or EXCLUDE), the constraint is renamed as well. There is no effect on the stored data.

SET TABLESPACE

This form changes the index’s tablespace to the specified tablespace and moves the data file(s) associated with the index to the new tablespace. To change the tablespace of an index, you must own the index and have CREATE privilege on the new tablespace. All indexes in the current database in a tablesipace can be moved by using the ALL IN TABLESPACE form, which will lock all indexes to be moved and then move each one. This form also supports OWNED BY, which will only move indexes owned by the roles specified. If the NOWAIT option is specified then the command will fail if it is unable to acquire all of the locks required immediately. Note that system catalogs will not be moved by this command, use ALTER DATABASE or explicit ALTER INDEX invocations instead if desired. See also CREATE TABLESPACE.

ATTACH PARTITION

Causes the named index to become attached to the altered index. The named index must be on a partition of the table containing the index being altered, and have an equivalent definition. An attached index cannot be dropped by itself, and will automatically be dropped if its parent index is dropped.

DEPENDS ON EXTENSION

This form marks the index as dependent on the extension, such that if the extension is dropped, the index will automatically be dropped as well.

SET ( storage_parameter = value [, ... ] )

This form changes one or more index-method-specific storage parameters for the index. See CREATE INDEX for details on the available parameters. Note that the index contents will not be
ALTER INDEX

modified immediately by this command; depending on the parameter you might need to rebuild
the index with REINDEX to get the desired effects.

RESET ( storage_parameter [, ... ] )

This form resets one or more index-method-specific storage parameters to their defaults. As with
SET, a REINDEX might be needed to update the index entirely.

ALTER [ COLUMN ] column_number SET STATISTICS integer

This form sets the per-column statistics-gathering target for subsequent ANALYZE operations,
though can be used only on index columns that are defined as an expression. Since expressions
lack a unique name, we refer to them using the ordinal number of the index column. The target
can be set in the range 0 to 10000; alternatively, set it to -1 to revert to using the system default
statistics target (default_statistics_target). For more information on the use of statistics by the
PostgreSQL query planner, refer to Section 14.2.

Parameters

IF EXISTS

Do not throw an error if the index does not exist. A notice is issued in this case.

column_number

The ordinal number refers to the ordinal (left-to-right) position of the index column.

name

The name (possibly schema-qualified) of an existing index to alter.

new_name

The new name for the index.

tablespace_name

The tablespace to which the index will be moved.

extension_name

The name of the extension that the index is to depend on.

storage_parameter

The name of an index-method-specific storage parameter.

value

The new value for an index-method-specific storage parameter. This might be a number or a word
depending on the parameter.

Notes

These operations are also possible using ALTER TABLE. ALTER INDEX is in fact just an alias for
the forms of ALTER TABLE that apply to indexes.

There was formerly an ALTER INDEX OWNER variant, but this is now ignored (with a warning). An
index cannot have an owner different from its table's owner. Changing the table's owner automatically
changes the index as well.
Changing any part of a system catalog index is not permitted.

**Examples**

To rename an existing index:

```sql
ALTER INDEX distributors RENAME TO suppliers;
```

To move an index to a different tablespace:

```sql
ALTER INDEX distributors SET TABLESPACE fasttablespace;
```

To change an index's fill factor (assuming that the index method supports it):

```sql
ALTER INDEX distributors SET (fillfactor = 75);
REINDEX INDEX distributors;
```

Set the statistics-gathering target for an expression index:

```sql
CREATE INDEX coord_idx ON measured (x, y, (z + t));
ALTER INDEX coord_idx ALTER COLUMN 3 SET STATISTICS 1000;
```

**Compatibility**

ALTER INDEX is a PostgreSQL extension.

**See Also**

CREATE INDEX, REINDEX
ALTER LANGUAGE

ALTER LANGUAGE — change the definition of a procedural language

Synopsis

ALTER [ PROCEDURAL ] LANGUAGE name RENAME TO new_name
ALTER [ PROCEDURAL ] LANGUAGE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER LANGUAGE changes the definition of a procedural language. The only functionality is to rename the language or assign a new owner. You must be superuser or owner of the language to use ALTER LANGUAGE.

Parameters

name

Name of a language

new_name

The new name of the language

new_owner

The new owner of the language

Compatibility

There is no ALTER LANGUAGE statement in the SQL standard.

See Also

CREATE LANGUAGE, DROP LANGUAGE
ALTER LARGE OBJECT
ALTER LARGE OBJECT — change the definition of a large object

Synopsis

ALTER LARGE OBJECT large_object_oid OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER LARGE OBJECT changes the definition of a large object. The only functionality is to assign a new owner. You must be superuser or owner of the large object to use ALTER LARGE OBJECT.

Parameters

large_object_oid
OID of the large object to be altered

new_owner
The new owner of the large object

Compatibility

There is no ALTER LARGE OBJECT statement in the SQL standard.

See Also
Chapter 35
ALTER MATERIALIZED VIEW

ALTER MATERIALIZED VIEW — change the definition of a materialized view

Synopsis

ALTER MATERIALIZED VIEW [ IF EXISTS ] name
   action [, ... ]
ALTER MATERIALIZED VIEW name
   DEPENDS ON EXTENSION extension_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
   RENAME [ COLUMN ] column_name TO new_column_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
   RENAME TO new_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
   SET SCHEMA new_schema
ALTER MATERIALIZED VIEW ALL IN TABLESPACE name [ OWNED BY role_name [, ... ] ]
   SET TABLESPACE new_tablespace [ NOWAIT ]

where action is one of:

   ALTER [ COLUMN ] column_name SET STATISTICS integer
   ALTER [ COLUMN ] column_name SET ( attribute_option = value [, ... ] )
   ALTER [ COLUMN ] column_name RESET ( attribute_option [, ... ] )
   ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL | EXTENDED | MAIN }
   CLUSTER ON index_name
   SET WITHOUT CLUSTER
   SET ( storage_parameter = value [, ... ] )
   RESET ( storage_parameter [, ... ] )
   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER MATERIALIZED VIEW changes various auxiliary properties of an existing materialized view.

You must own the materialized view to use ALTER MATERIALIZED VIEW. To change a materialized view's schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the materialized view's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the materialized view. However, a superuser can alter ownership of any view anyway.)

The DEPENDS ON EXTENSION form marks the materialized view as dependent on an extension, such that the materialized view will automatically be dropped if the extension is dropped.

The statement subforms and actions available for ALTER MATERIALIZED VIEW are a subset of those available for ALTER TABLE, and have the same meaning when used for materialized views. See the descriptions for ALTER TABLE for details.
ALTER MATERIALIZED VIEW

Parameters

name

The name (optionally schema-qualified) of an existing materialized view.

column_name

Name of a new or existing column.

extension_name

The name of the extension that the materialized view is to depend on.

new_column_name

New name for an existing column.

new_owner

The user name of the new owner of the materialized view.

new_name

The new name for the materialized view.

new_schema

The new schema for the materialized view.

Examples

To rename the materialized view foo to bar:

ALTER MATERIALIZED VIEW foo RENAME TO bar;

Compatibility

ALTER MATERIALIZED VIEW is a PostgreSQL extension.

See Also

CREATE MATERIALIZED VIEW, DROP MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
ALTER OPERATOR

ALTER OPERATOR — change the definition of an operator

Synopsis

```
ALTER OPERATOR name ( { left_type | NONE }, { right_type | NONE } )
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR name ( { left_type | NONE }, { right_type | NONE } )
    SET SCHEMA new_schema

ALTER OPERATOR name ( { left_type | NONE }, { right_type | NONE } )
    SET { { RESTRICT = { res_proc | NONE } |
               JOIN = { join_proc | NONE } }
         [, ... ] }
```

Description

ALTER OPERATOR changes the definition of an operator.

You must own the operator to use ALTER OPERATOR. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the operator's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the operator. However, a superuser can alter ownership of any operator anyway.)

Parameters

```
name

The name (optionally schema-qualified) of an existing operator.

left_type

The data type of the operator's left operand; write NONE if the operator has no left operand.

right_type

The data type of the operator's right operand; write NONE if the operator has no right operand.

new_owner

The new owner of the operator.

new_schema

The new schema for the operator.

res_proc

The restriction selectivity estimator function for this operator; write NONE to remove existing selectivity estimator.
```
join_proc

The join selectivity estimator function for this operator; write NONE to remove existing selectivity estimator.

Examples

Change the owner of a custom operator a @@ b for type text:

```
ALTER OPERATOR @@ (text, text) OWNER TO joe;
```

Change the restriction and join selectivity estimator functions of a custom operator a && b for type int[]:

```
ALTER OPERATOR && (_int4, _int4) SET (RESTRICT = _int_contsel, JOIN = _int_contjoinsel);
```

Compatibility

There is no ALTER OPERATOR statement in the SQL standard.

See Also

CREATE OPERATOR, DROP OPERATOR
ALTER OPERATOR CLASS

ALTER OPERATOR CLASS — change the definition of an operator class

Synopsis

ALTER OPERATOR CLASS name USING index_method
    RENAME TO new_name

ALTER OPERATOR CLASS name USING index_method
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR CLASS name USING index_method
    SET SCHEMA new_schema

Description

ALTER OPERATOR CLASS changes the definition of an operator class.

You must own the operator class to use ALTER OPERATOR CLASS. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the operator class's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the operator class. However, a superuser can alter ownership of any operator class anyway.)

Parameters

name
    The name (optionally schema-qualified) of an existing operator class.

index_method
    The name of the index method this operator class is for.

new_name
    The new name of the operator class.

new_owner
    The new owner of the operator class.

new_schema
    The new schema for the operator class.

Compatibility

There is no ALTER OPERATOR CLASS statement in the SQL standard.

See Also

CREATE OPERATOR CLASS, DROP OPERATOR CLASS, ALTER OPERATOR FAMILY
ALTER OPERATOR FAMILY

ALTER OPERATOR FAMILY — change the definition of an operator family

Synopsis

ALTER OPERATOR FAMILY name USING index_method ADD
{ OPERATOR strategy_number operator_name ( op_type, op_type )
   [ FOR SEARCH | FOR ORDER BY sort_family_name ]
   | FUNCTION support_number [ ( op_type [, , op_type ] ) ]
   function_name [ ( argument_type [, , ...] ) ]
 } [, ... ]

ALTER OPERATOR FAMILY name USING index_method DROP
{ OPERATOR strategy_number ( op_type [, , op_type ] )
   | FUNCTION support_number ( op_type [, , op_type ] )
 } [, ... ]

ALTER OPERATOR FAMILY name USING index_method
RENAME TO new_name

ALTER OPERATOR FAMILY name USING index_method
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR FAMILY name USING index_method
SET SCHEMA new_schema

Description

ALTER OPERATOR FAMILY changes the definition of an operator family. You can add operators and support functions to the family, remove them from the family, or change the family’s name or owner.

When operators and support functions are added to a family with ALTER OPERATOR FAMILY, they are not part of any specific operator class within the family, but are just “loose” within the family. This indicates that these operators and functions are compatible with the family’s semantics, but are not required for correct functioning of any specific index. (Operators and functions that are so required should be declared as part of an operator class, instead; see CREATE OPERATOR CLASS.) PostgreSQL will allow loose members of a family to be dropped from the family at any time, but members of an operator class cannot be dropped without dropping the whole class and any indexes that depend on it. Typically, single-data-type operators and functions are part of operator classes because they are needed to support an index on that specific data type, while cross-data-type operators and functions are made loose members of the family.

You must be a superuser to use ALTER OPERATOR FAMILY. (This restriction is made because an erroneous operator family definition could confuse or even crash the server.)

ALTER OPERATOR FAMILY does not presently check whether the operator family definition includes all the operators and functions required by the index method, nor whether the operators and functions form a self-consistent set. It is the user’s responsibility to define a valid operator family.

Refer to Section 38.15 for further information.
 Parameters

 name

 The name (optionally schema-qualified) of an existing operator family.

 index_method

 The name of the index method this operator family is for.

 strategy_number

 The index method's strategy number for an operator associated with the operator family.

 operator_name

 The name (optionally schema-qualified) of an operator associated with the operator family.

 op_type

 In an OPERATOR clause, the operand data type(s) of the operator, or NONE to signify a left-unary or right-unary operator. Unlike the comparable syntax in CREATE OPERATOR CLASS, the operand data types must always be specified.

 In an ADD FUNCTION clause, the operand data type(s) the function is intended to support, if different from the input data type(s) of the function. For B-tree comparison functions and hash functions it is not necessary to specify op_type since the function's input data type(s) are always the correct ones to use. For B-tree sort support functions and all functions in GiST, SP-GiST and GIN operator classes, it is necessary to specify the operand data type(s) the function is to be used with.

 In a DROP FUNCTION clause, the operand data type(s) the function is intended to support must be specified.

 sort_family_name

 The name (optionally schema-qualified) of an existing btree operator family that describes the sort ordering associated with an ordering operator.

 If neither FOR SEARCH nor FOR ORDER BY is specified, FOR SEARCH is the default.

 support_number

 The index method's support function number for a function associated with the operator family.

 function_name

 The name (optionally schema-qualified) of a function that is an index method support function for the operator family. If no argument list is specified, the name must be unique in its schema.

 argument_type

 The parameter data type(s) of the function.

 new_name

 The new name of the operator family.

 new_owner

 The new owner of the operator family.
new_schema

The new schema for the operator family.

The OPERATOR and FUNCTION clauses can appear in any order.

Notes

Notice that the DROP syntax only specifies the “slot” in the operator family, by strategy or support number and input data type(s). The name of the operator or function occupying the slot is not mentioned. Also, for DROP FUNCTION the type(s) to specify are the input data type(s) the function is intended to support; for GiST, SP-GiST and GIN indexes this might have nothing to do with the actual input argument types of the function.

Because the index machinery does not check access permissions on functions before using them, including a function or operator in an operator family is tantamount to granting public execute permission on it. This is usually not an issue for the sorts of functions that are useful in an operator family.

The operators should not be defined by SQL functions. A SQL function is likely to be inlined into the calling query, which will prevent the optimizer from recognizing that the query matches an index.

Before PostgreSQL 8.4, the OPERATOR clause could include a RECHECK option. This is no longer supported because whether an index operator is “lossy” is now determined on-the-fly at run time. This allows efficient handling of cases where an operator might or might not be lossy.

Examples

The following example command adds cross-data-type operators and support functions to an operator family that already contains B-tree operator classes for data types int4 and int2.

```
ALTER OPERATOR FAMILY integer_ops USING btree ADD

-- int4 vs int2
OPERATOR 1 < (int4, int2) ,
OPERATOR 2 <= (int4, int2) ,
OPERATOR 3 = (int4, int2) ,
OPERATOR 4 >= (int4, int2) ,
OPERATOR 5 > (int4, int2) ,
FUNCTION 1 btint42cmp(int4, int2) ,

-- int2 vs int4
OPERATOR 1 < (int2, int4) ,
OPERATOR 2 <= (int2, int4) ,
OPERATOR 3 = (int2, int4) ,
OPERATOR 4 >= (int2, int4) ,
OPERATOR 5 > (int2, int4) ,
FUNCTION 1 btint24cmp(int2, int4) ;
```

To remove these entries again:

```
ALTER OPERATOR FAMILY integer_ops USING btree DROP

-- int4 vs int2
OPERATOR 1 (int4, int2) ,
OPERATOR 2 (int4, int2) ,
OPERATOR 3 (int4, int2) ,
OPERATOR 4 (int4, int2) ,
```

1398
ALTER OPERATOR FAMILY

OPERATOR 5 (int4, int2),
FUNCTION 1 (int4, int2),

-- int2 vs int4
OPERATOR 1 (int2, int4),
OPERATOR 2 (int2, int4),
OPERATOR 3 (int2, int4),
OPERATOR 4 (int2, int4),
OPERATOR 5 (int2, int4),
FUNCTION 1 (int2, int4);

Compatibility

There is no ALTER OPERATOR FAMILY statement in the SQL standard.

See Also

CREATE OPERATOR FAMILY, DROP OPERATOR FAMILY, CREATE OPERATOR CLASS,
ALTER OPERATOR CLASS, DROP OPERATOR CLASS
ALTER POLICY

ALTER POLICY — change the definition of a row level security policy

Synopsis

ALTER POLICY name ON table_name RENAME TO new_name

ALTER POLICY name ON table_name
    [ TO { role_name | PUBLIC | CURRENT_USER | SESSION_USER } [, ...] ]
    [ USING { using_expression } ]
    [ WITH CHECK { check_expression } ]

Description

ALTER POLICY changes the definition of an existing row-level security policy. Note that ALTER POLICY only allows the set of roles to which the policy applies and the USING and WITH CHECK expressions to be modified. To change other properties of a policy, such as the command to which it applies or whether it is permissive or restrictive, the policy must be dropped and recreated.

To use ALTER POLICY, you must own the table that the policy applies to.

In the second form of ALTER POLICY, the role list, using_expression, and check_expression are replaced independently if specified. When one of those clauses is omitted, the corresponding part of the policy is unchanged.

Parameters

name

The name of an existing policy to alter.

table_name

The name (optionally schema-qualified) of the table that the policy is on.

new_name

The new name for the policy.

role_name

The role(s) to which the policy applies. Multiple roles can be specified at one time. To apply the policy to all roles, use PUBLIC.

using_expression

The USING expression for the policy. See CREATE POLICY for details.

check_expression

The WITH CHECK expression for the policy. See CREATE POLICY for details.

Compatibility

ALTER POLICY is a PostgreSQL extension.
See Also
CREATE POLICY, DROP POLICY
ALTER PROCEDURE — change the definition of a procedure

Synopsis

ALTER PROCEDURE name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] action [ ... ] [ RESTRICT ]
ALTER PROCEDURE name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ]
  RENAME TO new_name
ALTER PROCEDURE name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ]
  OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER PROCEDURE name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ]
  SET SCHEMA new_schema
ALTER PROCEDURE name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ]
  DEPENDS ON EXTENSION extension_name

where action is one of:

  [ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
SET configuration_parameter { TO | = } { value | DEFAULT }
SET configuration_parameter FROM CURRENT
RESET configuration_parameter
RESET ALL

Description

ALTER PROCEDURE changes the definition of a procedure.

You must own the procedure to use ALTER PROCEDURE. To change a procedure's schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the procedure's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the procedure. However, a superuser can alter ownership of any procedure anyway.)

Parameters

name

The name (optionally schema-qualified) of an existing procedure. If no argument list is specified, the name must be unique in its schema.

argmode

The mode of an argument: IN or VARIADIC. If omitted, the default is IN.
ALTER PROCEDURE

**argname**

The name of an argument. Note that ALTER PROCEDURE does not actually pay any attention to argument names, since only the argument data types are needed to determine the procedure's identity.

**argtype**

The data type(s) of the procedure's arguments (optionally schema-qualified), if any.

**new_name**

The new name of the procedure.

**new_owner**

The new owner of the procedure. Note that if the procedure is marked SECURITY DEFINER, it will subsequently execute as the new owner.

**new_schema**

The new schema for the procedure.

**extension_name**

The name of the extension that the procedure is to depend on.

[ EXTERNAL ] SECURITY INVOKER
[ EXTERNAL ] SECURITY DEFINER

Change whether the procedure is a security definer or not. The key word EXTERNAL is ignored for SQL conformance. See CREATE PROCEDURE for more information about this capability.

**configuration_parameter**

**value**

Add or change the assignment to be made to a configuration parameter when the procedure is called. If value is DEFAULT or, equivalently, RESET is used, the procedure-local setting is removed, so that the procedure executes with the value present in its environment. Use RESET ALL to clear all procedure-local settings. SET FROM CURRENT saves the value of the parameter that is current when ALTER PROCEDURE is executed as the value to be applied when the procedure is entered.

See SET and Chapter 19 for more information about allowed parameter names and values.

RESTRICT

Ignored for conformance with the SQL standard.

**Examples**

To rename the procedure insert_data with two arguments of type integer to insert_record:

```
ALTER PROCEDURE insert_data(integer, integer) RENAME TO
insert_record;
```

To change the owner of the procedure insert_data with two arguments of type integer to joe:
ALTER PROCEDURE insert_data(integer, integer) OWNER TO joe;

To change the schema of the procedure insert_data with two arguments of type integer to accounting:

ALTER PROCEDURE insert_data(integer, integer) SET SCHEMA accounting;

To mark the procedure insert_data(integer, integer) as being dependent on the extension myext:

ALTER PROCEDURE insert_data(integer, integer) DEPENDS ON EXTENSION myext;

To adjust the search path that is automatically set for a procedure:

ALTER PROCEDURE check_password(text) SET search_path = admin, pg_temp;

To disable automatic setting of search_path for a procedure:

ALTER PROCEDURE check_password(text) RESET search_path;

The procedure will now execute with whatever search path is used by its caller.

Compatibility

This statement is partially compatible with the ALTER PROCEDURE statement in the SQL standard. The standard allows more properties of a procedure to be modified, but does not provide the ability to rename a procedure, make a procedure a security definer, attach configuration parameter values to a procedure, or change the owner, schema, or volatility of a procedure. The standard also requires the RESTRICT key word, which is optional in PostgreSQL.

See Also

CREATE PROCEDURE, DROP PROCEDURE, ALTER FUNCTION, ALTER ROUTINE
ALTER PUBLICATION

ALTER PUBLICATION — change the definition of a publication

Synopsis

ALTER PUBLICATION name ADD TABLE [ ONLY ] table_name [ * ] [, ...]
ALTER PUBLICATION name SET TABLE [ ONLY ] table_name [ * ] [, ...]
ALTER PUBLICATION name DROP TABLE [ ONLY ] table_name [ * ] [, ...]
ALTER PUBLICATION name SET ( publication_parameter [= value] [, ... ] )
ALTER PUBLICATION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER PUBLICATION name RENAME TO new_name

Description

The command ALTER PUBLICATION can change the attributes of a publication.

The first three variants change which tables are part of the publication. The SET TABLE clause will replace the list of tables in the publication with the specified one. The ADD TABLE and DROP TABLE clauses will add and remove one or more tables from the publication. Note that adding tables to a publication that is already subscribed to will require a ALTER SUBSCRIPTION ... REFRESH PUBLICATION action on the subscribing side in order to become effective.

The fourth variant of this command listed in the synopsis can change all of the publication properties specified in CREATE PUBLICATION. Properties not mentioned in the command retain their previous settings.

The remaining variants change the owner and the name of the publication.

You must own the publication to use ALTER PUBLICATION. To alter the owner, you must also be a direct or indirect member of the new owning role. The new owner must have CREATE privilege on the database. Also, the new owner of a FOR ALL TABLES publication must be a superuser. However, a superuser can change the ownership of a publication while circumventing these restrictions.

Parameters

name

The name of an existing publication whose definition is to be altered.

table_name

Name of an existing table. If ONLY is specified before the table name, only that table is affected. If ONLY is not specified, the table and all its descendant tables (if any) are affected. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

SET ( publication_parameter [= value] [, ... ] )

This clause alters publication parameters originally set by CREATE PUBLICATION. See there for more information.

new_owner

The user name of the new owner of the publication.
new_name

The new name for the publication.

Examples

Change the publication to publish only deletes and updates:

```
ALTER PUBLICATION noinsert SET (publish = 'update, delete');
```

Add some tables to the publication:

```
ALTER PUBLICATION mypublication ADD TABLE users, departments;
```

Compatibility

**ALTER PUBLICATION** is a PostgreSQL extension.

See Also

CREATE PUBLICATION, DROP PUBLICATION, CREATE SUBSCRIPTION, ALTER SUBSCRIPTION
ALTER ROLE

ALTER ROLE — change a database role

Synopsis

ALTER ROLE role_specification [ WITH ] option [ ... ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICATION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT conlimit
- [ ENCRYPTED ] PASSWORD 'password'
- VALID UNTIL 'timestamp'

ALTER ROLE name RENAME TO new_name

ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ]
SET configuration_parameter { TO | = } { value | DEFAULT }

ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ]
SET configuration_parameter FROM CURRENT

ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ]
RESET configuration_parameter

ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ]
RESET ALL

where role_specification can be:

- role_name
- CURRENT_USER
- SESSION_USER

Description

ALTER ROLE changes the attributes of a PostgreSQL role.

The first variant of this command listed in the synopsis can change many of the role attributes that can be specified in CREATE ROLE. (All the possible attributes are covered, except that there are no options for adding or removing memberships; use GRANT and REVOKE for that.) Attributes not mentioned in the command retain their previous settings. Database superusers can change any of these settings for any role. Roles having CREATEROLE privilege can change any of these settings, but only for non-superuser and non-replication roles. Ordinary roles can only change their own password.

The second variant changes the name of the role. Database superusers can rename any role. Roles having CREATEROLE privilege can rename non-superuser roles. The current session user cannot be renamed. (Connect as a different user if you need to do that.) Because MD5-encrypted passwords use the role name as cryptographic salt, renaming a role clears its password if the password is MD5-encrypted.
The remaining variants change a role’s session default for a configuration variable, either for all databases or, when the `IN DATABASE` clause is specified, only for sessions in the named database. If `ALL` is specified instead of a role name, this changes the setting for all roles. Using `ALL` with `IN DATABASE` is effectively the same as using the command `ALTER DATABASE ... SET ....`

Whenever the role subsequently starts a new session, the specified value becomes the session default, overriding whatever setting is present in `postgresql.conf` or has been received from the `postgres` command line. This only happens at login time; executing `SET ROLE` or `SET SESSION AUTHORIZATION` does not cause new configuration values to be set. Settings set for all databases are overridden by database-specific settings attached to a role. Settings for specific databases or specific roles override settings for all roles.

Superusers can change anyone's session defaults. Roles having `CREATEROLE` privilege can change defaults for non-superuser roles. Ordinary roles can only set defaults for themselves. Certain configuration variables cannot be set this way, or can only be set if a superuser issues the command. Only superusers can change a setting for all roles in all databases.

**Parameters**

`name`

The name of the role whose attributes are to be altered.

`CURRENT_USER`

Alter the current user instead of an explicitly identified role.

`SESSION_USER`

Alter the current session user instead of an explicitly identified role.

`SUPERUSER`

`NOSUPERUSER`

`CREATEDB`

`NOCREATEDB`

`CREATEROLE`

`NOCREATEROLE`

`INHERIT`

`NOINHERIT`

`LOGIN`

`NOLOGIN`

`REPLICATION`

`NOREPLICATION`

`BYPASSRLS`

`NOBYPASSRLS`

`CONNECTION LIMIT connlimit`

`[ ENCRYPTED ] PASSWORD password`

`VALID UNTIL 'timestamp'`

These clauses alter attributes originally set by `CREATE ROLE`. For more information, see the `CREATE ROLE` reference page.

`new_name`

The new name of the role.

`database_name`

The name of the database the configuration variable should be set in.
ALTER ROLE

classification_parameter
value

Set this role's session default for the specified configuration parameter to the given value. If value is DEFAULT or, equivalently, RESET is used, the role-specific variable setting is removed, so the role will inherit the system-wide default setting in new sessions. Use RESET ALL to clear all role-specific settings. SET FROM CURRENT saves the session's current value of the parameter as the role-specific value. If IN DATABASE is specified, the configuration parameter is set or removed for the given role and database only.

Role-specific variable settings take effect only at login; SET ROLE and SET SESSION AUTHORIZATION do not process role-specific variable settings.

See SET and Chapter 19 for more information about allowed parameter names and values.

Notes

Use CREATE ROLE to add new roles, and DROP ROLE to remove a role.

ALTER ROLE cannot change a role's memberships. Use GRANT and REVOKE to do that.

Caution must be exercised when specifying an unencrypted password with this command. The password will be transmitted to the server in cleartext, and it might also be logged in the client's command history or the server log. psql contains a command \password that can be used to change a role's password without exposing the cleartext password.

It is also possible to tie a session default to a specific database rather than to a role; see ALTER DATABASE. If there is a conflict, database-role-specific settings override role-specific ones, which in turn override database-specific ones.

Examples

Change a role's password:

```
ALTER ROLE davide WITH PASSWORD 'hu8jmn3';
```

Remove a role's password:

```
ALTER ROLE davide WITH PASSWORD NULL;
```

Change a password expiration date, specifying that the password should expire at midday on 4th May 2015 using the time zone which is one hour ahead of UTC:

```
ALTER ROLE chris VALID UNTIL 'May 4 12:00:00 2015 +1';
```

Make a password valid forever:

```
ALTER ROLE fred VALID UNTIL 'infinity';
```

Give a role the ability to create other roles and new databases:

```
ALTER ROLE miriam CREATEROLE CREATEDB;
```

Give a role a non-default setting of the maintenance_work_mem parameter:
ALTER ROLE worker_bee SET maintenance_work_mem = 100000;

Give a role a non-default, database-specific setting of the client_min_messages parameter:

ALTER ROLE fred IN DATABASE devel SET client_min_messages = DEBUG;

**Compatibility**

The `ALTER ROLE` statement is a PostgreSQL extension.

**See Also**

CREATE ROLE, DROP ROLE, ALTER DATABASE, SET
ALTER ROUTINE

ALTER ROUTINE — change the definition of a routine

Synopsis

```
ALTER ROUTINE name [ ( [ [ argmode ] [ argname ] argtype
{, ...] ] ) ]
   action [ ... ] [ RESTRICT ]
ALTER ROUTINE name [ ( [ [ argmode ] [ argname ] argtype
{, ...] ] ) ]
   RENAME TO new_name
ALTER ROUTINE name [ ( [ [ argmode ] [ argname ] argtype
{, ...] ] ) ]
   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER ROUTINE name [ ( [ [ argmode ] [ argname ] argtype
{, ...] ] ) ]
   SET SCHEMA new_schema
ALTER ROUTINE name [ ( [ [ argmode ] [ argname ] argtype
{, ...] ] ) ]
   DEPENDS ON EXTENSION extension_name
```

Where `action` is one of:

```
IMMUTABLE | STABLE | VOLATILE | [ NOT ] LEAKPROOF
[ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
PARALLEL { UNSAFE | RESTRICTED | SAFE }
COST execution_cost
ROWS result_rows
SET configuration_parameter { TO | = } { value | DEFAULT }
SET configuration_parameter FROM CURRENT
RESET configuration_parameter
RESET ALL
```

Description

ALTER ROUTINE changes the definition of a routine, which can be an aggregate function, a normal function, or a procedure. See under ALTER AGGREGATE, ALTER FUNCTION, and ALTER PROCEDURE for the description of the parameters, more examples, and further details.

Examples

To rename the routine `foo` for type integer to `foobar`:

```
ALTER ROUTINE foo(integer) RENAME TO foobar;
```

This command will work independent of whether `foo` is an aggregate, function, or procedure.

Compatibility

This statement is partially compatible with the ALTER ROUTINE statement in the SQL standard. See under ALTER FUNCTION and ALTER PROCEDURE for more details. Allowing routine names to refer to aggregate functions is a PostgreSQL extension.
See Also

ALTER AGGREGATE, ALTER FUNCTION, ALTER PROCEDURE, DROP ROUTINE

Note that there is no CREATE ROUTINE command.
ALTER RULE

ALTER RULE — change the definition of a rule

Synopsis

ALTER RULE name ON table_name RENAME TO new_name

Description

ALTER RULE changes properties of an existing rule. Currently, the only available action is to change the rule's name.

To use ALTER RULE, you must own the table or view that the rule applies to.

Parameters

*name*

The name of an existing rule to alter.

*table_name*

The name (optionally schema-qualified) of the table or view that the rule applies to.

*new_name*

The new name for the rule.

Examples

To rename an existing rule:

```
ALTER RULE notify_all ON emp RENAME TO notify_me;
```

Compatibility

ALTER RULE is a PostgreSQL language extension, as is the entire query rewrite system.

See Also

CREATE RULE, DROP RULE
ALTER SCHEMA

ALTER SCHEMA — change the definition of a schema

Synopsis

ALTER SCHEMA name RENAME TO new_name
ALTER SCHEMA name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER SCHEMA changes the definition of a schema.

You must own the schema to use ALTER SCHEMA. To rename a schema you must also have the CREATE privilege for the database. To alter the owner, you must also be a direct or indirect member of the new owning role, and you must have the CREATE privilege for the database. (Note that superusers have all these privileges automatically.)

Parameters

name

The name of an existing schema.

new_name

The new name of the schema. The new name cannot begin with pg_, as such names are reserved for system schemas.

new_owner

The new owner of the schema.

Compatibility

There is no ALTER SCHEMA statement in the SQL standard.

See Also

CREATE SCHEMA, DROP SCHEMA
ALTER SEQUENCE

ALTER SEQUENCE — change the definition of a sequence generator

Synopsis

ALTER SEQUENCE [ IF EXISTS ] name
   [ AS data_type ]
   [ INCREMENT [ BY ] increment ]
   [ MINVALUE minvalue | NO MINVALUE ] [ MAXVALUE maxvalue | NO MAXVALUE ]
   [ START [ WITH ] start ]
   [ RESTART [ [ WITH ] restart ] ]
   [ CACHE cache ] [ [ NO ] CYCLE ]
   [ OWNED BY { table_name.column_name | NONE } ]
ALTER SEQUENCE [ IF EXISTS ] name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER SEQUENCE [ IF EXISTS ] name RENAME TO new_name
ALTER SEQUENCE [ IF EXISTS ] name SET SCHEMA new_schema

Description

ALTER SEQUENCE changes the parameters of an existing sequence generator. Any parameters not specifically set in the ALTER SEQUENCE command retain their prior settings.

You must own the sequence to use ALTER SEQUENCE. To change a sequence's schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the sequence's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the sequence. However, a superuser can alter ownership of any sequence anyway.)

Parameters

 name

The name (optionally schema-qualified) of a sequence to be altered.

IF EXISTS

Do not throw an error if the sequence does not exist. A notice is issued in this case.

data_type

The optional clause AS data_type changes the data type of the sequence. Valid types are smallint, integer, and bigint.

Changing the data type automatically changes the minimum and maximum values of the sequence if and only if the previous minimum and maximum values were the minimum or maximum value of the old data type (in other words, if the sequence had been created using NO MINVALUE or NO MAXVALUE, implicitly or explicitly). Otherwise, the minimum and maximum values are preserved, unless new values are given as part of the same command. If the minimum and maximum values do not fit into the new data type, an error will be generated.
increment

The clause INCREMENT BY increment is optional. A positive value will make an ascending sequence, a negative one a descending sequence. If unspecified, the old increment value will be maintained.

minvalue

NO MINVALUE

The optional clause MINVALUE minvalue determines the minimum value a sequence can generate. If NO MINVALUE is specified, the defaults of 1 and the minimum value of the data type for ascending and descending sequences, respectively, will be used. If neither option is specified, the current minimum value will be maintained.

maxvalue

NO MAXVALUE

The optional clause MAXVALUE maxvalue determines the maximum value for the sequence. If NO MAXVALUE is specified, the defaults of the maximum value of the data type and -1 for ascending and descending sequences, respectively, will be used. If neither option is specified, the current maximum value will be maintained.

start

The optional clause START WITH start changes the recorded start value of the sequence. This has no effect on the current sequence value; it simply sets the value that future ALTER SEQUENCE RESTART commands will use.

restart

The optional clause RESTART [ WITH restart ] changes the current value of the sequence. This is similar to calling the setval function with is_called = false: the specified value will be returned by the next call of nextval. Writing RESTART with no restart value is equivalent to supplying the start value that was recorded by CREATE SEQUENCE or last set by ALTER SEQUENCE START WITH.

In contrast to a setval call, a RESTART operation on a sequence is transactional and blocks concurrent transactions from obtaining numbers from the same sequence. If that's not the desired mode of operation, setval should be used.

cache

The clause CACHE cache enables sequence numbers to be preallocated and stored in memory for faster access. The minimum value is 1 (only one value can be generated at a time, i.e., no cache). If unspecified, the old cache value will be maintained.

CYCLE

The optional CYCLE key word can be used to enable the sequence to wrap around when the maxvalue or minvalue has been reached by an ascending or descending sequence respectively. If the limit is reached, the next number generated will be the minvalue or maxvalue, respectively.

NO CYCLE

If the optional NO CYCLE key word is specified, any calls to nextval after the sequence has reached its maximum value will return an error. If neither CYCLE or NO CYCLE are specified, the old cycle behavior will be maintained.

OWNED BY table_name.column_name
OWNED BY NONE

The OWNED BY option causes the sequence to be associated with a specific table column, such that if that column (or its whole table) is dropped, the sequence will be automatically dropped as well.
If specified, this association replaces any previously specified association for the sequence. The specified table must have the same owner and be in the same schema as the sequence. Specifying OWNED BY NONE removes any existing association, making the sequence “free-standing”.

new_owner

The user name of the new owner of the sequence.

new_name

The new name for the sequence.

new_schema

The new schema for the sequence.

Notes

ALTER SEQUENCE will not immediately affect nextval results in backends, other than the current one, that have preallocated (cached) sequence values. They will use up all cached values prior to noticing the changed sequence generation parameters. The current backend will be affected immediately.

ALTER SEQUENCE does not affect the currval status for the sequence. (Before PostgreSQL 8.3, it sometimes did.)

ALTER SEQUENCE blocks concurrent nextval, currval, lastval, and setval calls.

For historical reasons, ALTER TABLE can be used with sequences too; but the only variants of ALTER TABLE that are allowed with sequences are equivalent to the forms shown above.

Examples

Restart a sequence called serial, at 105:

ALTER SEQUENCE serial RESTART WITH 105;

Compatibility

ALTER SEQUENCE conforms to the SQL standard, except for the AS, START WITH, OWNED BY, OWNER TO, RENAME TO, and SET SCHEMA clauses, which are PostgreSQL extensions.

See Also

CREATE SEQUENCE, DROP SEQUENCE
ALTER SERVER

ALTER SERVER — change the definition of a foreign server

Synopsis

```
ALTER SERVER name [ VERSION 'new_version' ]
   [ OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] ) ]
ALTER SERVER name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER SERVER name RENAME TO new_name
```

Description

ALTER SERVER changes the definition of a foreign server. The first form changes the server version string or the generic options of the server (at least one clause is required). The second form changes the owner of the server.

To alter the server you must be the owner of the server. Additionally to alter the owner, you must own the server and also be a direct or indirect member of the new owning role, and you must have USAGE privilege on the server's foreign-data wrapper. (Note that superusers satisfy all these criteria automatically.)

Parameters

- `name`
  The name of an existing server.

- `new_version`
  New server version.

- `OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )`
  Change options for the server. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Option names must be unique; names and values are also validated using the server's foreign-data wrapper library.

- `new_owner`
  The user name of the new owner of the foreign server.

- `new_name`
  The new name for the foreign server.

Examples

Alter server `foo`, add connection options:

```
ALTER SERVER foo OPTIONS (host 'foo', dbname 'foodb');
```

Alter server `foo`, change version, change host option:
ALTER SERVER foo VERSION '8.4' OPTIONS (SET host 'baz');

Compatibility

ALTER SERVER conforms to ISO/IEC 9075-9 (SQL/MED). The OWNER TO and RENAME forms are PostgreSQL extensions.

See Also
CREATE SERVER, DROP SERVER
ALTER STATISTICS

ALTER STATISTICS — change the definition of an extended statistics object

Synopsis

ALTER STATISTICS name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER STATISTICS name RENAME TO new_name
ALTER STATISTICS name SET SCHEMA new_schema

Description

ALTER STATISTICS changes the parameters of an existing extended statistics object. Any parameters not specifically set in the ALTER STATISTICS command retain their prior settings.

You must own the statistics object to use ALTER STATISTICS. To change a statistics object’s schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the statistics object’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the statistics object. However, a superuser can alter ownership of any statistics object anyway.)

Parameters

name

The name (optionally schema-qualified) of the statistics object to be altered.

new_owner

The user name of the new owner of the statistics object.

new_name

The new name for the statistics object.

new_schema

The new schema for the statistics object.

Compatibility

There is no ALTER STATISTICS command in the SQL standard.

See Also

CREATE STATISTICS, DROP STATISTICS
ALTER SUBSCRIPTION

ALTER SUBSCRIPTION — change the definition of a subscription

Synopsis

ALTER SUBSCRIPTION name CONNECTION 'conninfo'
ALTER SUBSCRIPTION name SET PUBLICATION publication_name [, ...]
    [ WITH ( set_publication_option [= value] [, ... ] ) ]
ALTER SUBSCRIPTION name REFRESH PUBLICATION [ WITH ( refresh_option
    [= value] [, ... ] ) ]
ALTER SUBSCRIPTION name ENABLE
ALTER SUBSCRIPTION name DISABLE
ALTER SUBSCRIPTION name SET ( subscription_parameter [= value]
    [, ... ] )
ALTER SUBSCRIPTION name OWNER TO { new_owner | CURRENT_USER |
    SESSION_USER } ALTER SUBSCRIPTION name RENAME TO new_name

Description

ALTER SUBSCRIPTION can change most of the subscription properties that can be specified in CREATE SUBSCRIPTION.

You must own the subscription to use ALTER SUBSCRIPTION. To alter the owner, you must also be a direct or indirect member of the new owning role. The new owner has to be a superuser. (Currently, all subscription owners must be superusers, so the owner checks will be bypassed in practice. But this might change in the future.)

Parameters

name

The name of a subscription whose properties are to be altered.

CONNECTION 'conninfo'

This clause alters the connection property originally set by CREATE SUBSCRIPTION. See there for more information.

SET PUBLICATION publication_name

Changes list of subscribed publications. See CREATE SUBSCRIPTION for more information. By default this command will also act like REFRESH PUBLICATION.

set_publication_option specifies additional options for this operation. The supported options are:

refresh(boolean)

When false, the command will not try to refresh table information. REFRESH PUBLICATION should then be executed separately. The default is true.

Additionally, refresh options as described under REFRESH PUBLICATION may be specified.
ALTER SUBSCRIPTION

REFRESH PUBLICATION

Fetch missing table information from publisher. This will start replication of tables that were added to the subscribed-to publications since the last invocation of REFRESH PUBLICATION or since CREATE SUBSCRIPTION.

refresh_option specifies additional options for the refresh operation. The supported options are:

copy_data (boolean)

Specifies whether the existing data in the publications that are being subscribed to should be copied once the replication starts. The default is true.

ENABLE

Enables the previously disabled subscription, starting the logical replication worker at the end of transaction.

DISABLE

Disables the running subscription, stopping the logical replication worker at the end of transaction.

SET ( subscription_parameter [= value] [, ... ] )

This clause alters parameters originally set by CREATE SUBSCRIPTION. See there for more information. The allowed options are slot_name and synchronous_commit

new_owner

The user name of the new owner of the subscription.

new_name

The new name for the subscription.

Examples

Change the publication subscribed by a subscription to insert_only:

ALTER SUBSCRIPTION mysub SET PUBLICATION insert_only;

Disable (stop) the subscription:

ALTER SUBSCRIPTION mysub DISABLE;

Compatibility

ALTER SUBSCRIPTION is a PostgreSQL extension.

See Also

CREATE SUBSCRIPTION, DROP SUBSCRIPTION, CREATE PUBLICATION, ALTER PUBLICATION
ALTER SYSTEM

ALTER SYSTEM — change a server configuration parameter

Synopsis

ALTER SYSTEM SET configuration_parameter ( TO | = ) ( value | 'value' | DEFAULT )

ALTER SYSTEM RESET configuration_parameter
ALTER SYSTEM RESET ALL

Description

ALTER SYSTEM is used for changing server configuration parameters across the entire database cluster. It can be more convenient than the traditional method of manually editing the postgresql.conf file. ALTER SYSTEM writes the given parameter setting to the postgresql.auto.conf file, which is read in addition to postgresql.conf. Setting a parameter to DEFAULT, or using the RESET variant, removes that configuration entry from the postgresql.auto.conf file. Use RESET ALL to remove all such configuration entries.

Values set with ALTER SYSTEM will be effective after the next server configuration reload, or after the next server restart in the case of parameters that can only be changed at server start. A server configuration reload can be commanded by calling the SQL function pg_reload_conf(), running pg_ctl reload, or sending a SIGHUP signal to the main server process.

Only superusers can use ALTER SYSTEM. Also, since this command acts directly on the file system and cannot be rolled back, it is not allowed inside a transaction block or function.

Parameters

configuration_parameter

Name of a settable configuration parameter. Available parameters are documented in Chapter 19.

value

New value of the parameter. Values can be specified as string constants, identifiers, numbers, or comma-separated lists of these, as appropriate for the particular parameter. DEFAULT can be written to specify removing the parameter and its value from postgresql.auto.conf.

Notes

This command can’t be used to set data_directory, nor parameters that are not allowed in postgresql.conf (e.g., preset options).

See Section 19.1 for other ways to set the parameters.

Examples

Set the wal_level:

ALTER SYSTEM SET wal_level = replica;
ALTER SYSTEM

Undo that, restoring whatever setting was effective in `postgresql.conf`:

```
ALTER SYSTEM RESET wal_level;
```

**Compatibility**

The `ALTER SYSTEM` statement is a PostgreSQL extension.

**See Also**

`SET`, `SHOW`
ALTER TABLE

ALTER TABLE — change the definition of a table

Synopsis

ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
  action [, ... ]
ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
  RENAME [ COLUMN ] column_name TO new_column_name
ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
  RENAME CONSTRAINT constraint_name TO new_constraint_name
ALTER TABLE [ IF EXISTS ] name
  RENAME TO new_name
ALTER TABLE [ IF EXISTS ] name
  SET SCHEMA new_schema
ALTER TABLE ALL IN TABLESPACE name [ OWNED BY role_name [, ... ] ]
  SET TABLESPACE new_tablespace [ NOWAIT ]
ALTER TABLE [ IF EXISTS ] name
  ATTACH PARTITION partition_name { FOR
  VALUES partition_bound_spec | DEFAULT }
ALTER TABLE [ IF EXISTS ] name
  DETACH PARTITION partition_name

where action is one of:

  ADD [ COLUMN ] [ IF NOT EXISTS ] column_name data_type
  [ COLLATE collation ] [ column_constraint [ ... ] ]
  DROP [ COLUMN ] [ IF EXISTS ] column_name [ RESTRICT ]
  CASCADE ]
ALTER [ COLUMN ] column_name [ SET DATA ] TYPE data_type
  [ COLLATE collation ] [ USING expression ]
ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER [ COLUMN ] column_name ( SET | DROP ) NOT NULL
ALTER [ COLUMN ] column_name ADD GENERATED ( ALWAYS | BY
  DEFAULT ) AS IDENTITY [ ( sequence_options ) ]
ALTER [ COLUMN ] column_name ( SET GENERATED { ALWAYS | BY
  DEFAULT } ) SET sequence_option [ RESTART [ [ WITH ] restart ] }
[...]
ALTER [ COLUMN ] column_name DROP IDENTITY [ IF EXISTS ]
ALTER [ COLUMN ] column_name SET STATISTICS integer
ALTER [ COLUMN ] column_name SET ( attribute_option = value
  [ , ... ] )
ALTER [ COLUMN ] column_name RESET ( attribute_option
  [ , ... ] )
ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL |
  EXTENDED | MAIN }
ADD table_constraint [ NOT VALID ]
ADD table_constraint_using_index
ALTER CONSTRAINT constraint_name [ DEFERRABLE | NOT
DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]
VALIDATE CONSTRAINT constraint_name
DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT ]
CASCADE ]
ALTER TABLE

DISABLE TRIGGER [ \trigger_name | ALL | USER ]
ENABLE TRIGGER [ \trigger_name | ALL | USER ]
ENABLE REPLICA TRIGGER \trigger_name
ENABLE ALWAYS TRIGGER \trigger_name
DISABLE RULE rewrite_rule_name
ENABLE RULE rewrite_rule_name
ENABLE REPLICA RULE rewrite_rule_name
ENABLE ALWAYS RULE rewrite_rule_name
DISABLE ROW LEVEL SECURITY
ENABLE ROW LEVEL SECURITY
FORCE ROW LEVEL SECURITY
NO FORCE ROW LEVEL SECURITY
CLUSTER ON index_name
SET WITHOUT CLUSTER
SET WITH OIDS
SET WITHOUT OIDS
SET TABLESPACE new_tablespace
SET ( LOGGED | UNLOGGED )
SET ( storage_parameter = value [, ... ] )
RESET ( storage_parameter [, ... ] )
INHERIT parent_table
NO INHERIT parent_table
OF type_name
NOT OF
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
REPLICA IDENTITY ( DEFAULT | USING INDEX index_name | FULL | NOTHING )
and partition_bound_spec is:

IN ( ( numeric_literal | string_literal | TRUE | FALSE | NULL ) [, ... ] )
FROM ( ( numeric_literal | string_literal | TRUE | FALSE | MINVALUE | MAXVALUE ) [, ... ] )
TO ( ( numeric_literal | string_literal | TRUE | FALSE | MINVALUE | MAXVALUE ) [, ... ] )
WITH ( MODULUS numeric_literal, REMAINDER numeric_literal )

and column_constraint is:

[ CONSTRAINT constraint_name ]
| NOT NULL |
| NULL |
| CHECK ( expression ) [ NO INHERIT ] |
| DEFAULT default_expr |
| GENERATED { ALWAYS | BY DEFAULT } AS IDENTITY |
| ( sequence_options ) |
| UNIQUE index_parameters |
| PRIMARY KEY index_parameters |
| REFERENCES reftable ( ( refcolumn ) ) [ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ]
| [ ON DELETE action ] [ ON UPDATE action ] |
| [ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

and table_constraint is:

[ CONSTRAINT constraint_name ]
ALTER TABLE

{ CHECK ( expression ) | NO INHERIT | UNIQUE ( column_name [, ... ] ) index_parameters | PRIMARY KEY ( column_name [, ... ] ) index_parameters | EXCLUDE [ USING index_method ] { exclude_element WITH operator [, ... ] index_parameters | WHERE ( predicate ) } | FOREIGN KEY ( column_name [, ... ] ) REFERENCES reftable | ( refcolumn [, ... ] ) ] } | [ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ] | [ ON DELETE action | [ ON UPDATE action ] } | [ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

and table_constraint_using_index is:

[ CONSTRAINT constraint_name ]
{ UNIQUE | PRIMARY KEY } USING INDEX index_name | [ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

index_parameters in UNIQUE, PRIMARY KEY, and EXCLUDE constraints are:

[ INCLUDE ( column_name [, ... ] ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ USING INDEX TABLESPACE tablespace_name ]

exclude_element in an EXCLUDE constraint is:

{ column_name | ( expression ) } [ opclass ] [ ASC | DESC ] [ NULLS { FIRST | LAST } ]

Description

ALTER TABLE changes the definition of an existing table. There are several subforms described below. Note that the lock level required may differ for each subform. An ACCESS EXCLUSIVE lock is held unless explicitly noted. When multiple subcommands are listed, the lock held will be the strictest one required from any subcommand.

ADD COLUMN [ IF NOT EXISTS ]

This form adds a new column to the table, using the same syntax as CREATE TABLE. If IF NOT EXISTS is specified and a column already exists with this name, no error is thrown.

DROP COLUMN [ IF EXISTS ]

This form drops a column from a table. Indexes and table constraints involving the column will be automatically dropped as well. Multivariate statistics referencing the dropped column will also be removed if the removal of the column would cause the statistics to contain data for only a single column. You will need to say CASCADE if anything outside the table depends on the column, for example, foreign key references or views. If IF EXISTS is specified and the column does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of a column of a table. Indexes and simple table constraints involving the column will be automatically converted to use the new column type by reparseing the originally supplied expression. The optional COLLATE clause specifies a collation for the new column; if omitted, the collation is the default for the new column type. The optional USING clause specifies how to compute the new column value from the old; if omitted, the default conversion is the same.
as an assignment cast from old data type to new. A USING clause must be provided if there is no implicit or assignment cast from old to new type.

SET/DROP DEFAULT

These forms set or remove the default value for a column. Default values only apply in subsequent INSERT or UPDATE commands; they do not cause rows already in the table to change.

SET/DROP NOT NULL

These forms change whether a column is marked to allow null values or to reject null values. You can only use SET NOT NULL when the column contains no null values.

If this table is a partition, one cannot perform DROP NOT NULL on a column if it is marked NOT NULL in the parent table. To drop the NOT NULL constraint from all the partitions, perform DROP NOT NULL on the parent table. Even if there is no NOT NULL constraint on the parent, such a constraint can still be added to individual partitions, if desired; that is, the children can disallow nulls even if the parent allows them, but not the other way around.

ADD GENERATED { ALWAYS | BY DEFAULT } AS IDENTITY
SET GENERATED { ALWAYS | BY DEFAULT }
DROP IDENTITY [ IF EXISTS ]

These forms change whether a column is an identity column or change the generation attribute of an existing identity column. See CREATE TABLE for details.

If DROP IDENTITY IF EXISTS is specified and the column is not an identity column, no error is thrown. In this case a notice is issued instead.

SET sequence_option
RESTART

These forms alter the sequence that underlies an existing identity column. sequence_option is an option supported by ALTER SEQUENCE such as INCREMENT BY.

SET STATISTICS

This form sets the per-column statistics-gathering target for subsequent ANALYZE operations. The target can be set in the range 0 to 10000; alternatively, set it to -1 to revert to using the system default statistics target (default_statistics_target). For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.

SET STATISTICS acquires a SHARE UPDATE EXCLUSIVE lock.

SET ( attribute_option = value [, ... ] )
RESET ( attribute_option [, ... ] )

This form sets or resets per-attribute options. Currently, the only defined per-attribute options are n_distinct and n_distinct_inherited, which override the number-of-distinct-values estimates made by subsequent ANALYZE operations. n_distinct affects the statistics for the table itself, while n_distinct_inherited affects the statistics gathered for the table plus its inheritance children. When set to a positive value, ANALYZE will assume that the column contains exactly the specified number of distinct nonnull values. When set to a negative value, which must be greater than or equal to -1, ANALYZE will assume that the number of distinct nonnull values in the column is linear in the size of the table; the exact count is to be computed by multiplying the estimated table size by the absolute value of the given number. For example, a value of -1 implies that all values in the column are distinct, while a value of -0.5 implies that each value appears twice on the average. This can be useful when the size of the table changes over time, since the multiplication by the number of rows in the table is not performed until query planning time. Specify a value of 0 to revert to estimating the number of distinct values normally. For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.
ALTER TABLE

Changing per-attribute options acquires a SHARE UPDATE EXCLUSIVE lock.

SET STORAGE

This form sets the storage mode for a column. This controls whether this column is held inline or in a secondary TOAST table, and whether the data should be compressed or not. PLAIN must be used for fixed-length values such as integer and is inline, uncompressed. MAIN is for inline, compressible data. EXTERNAL is for external, uncompressed data, and EXTENDED is for external, compressed data. EXTENDED is the default for most data types that support non-PLAIN storage. Use of EXTERNAL will make substring operations on very large text and bytea values run faster, at the penalty of increased storage space. Note that SET STORAGE doesn't itself change anything in the table; it just sets the strategy to be pursued during future table updates. See Section 68.2 for more information.

ADD table_constraint [ NOT VALID ]

This form adds a new constraint to a table using the same syntax as CREATE TABLE, plus the option NOT VALID, which is currently only allowed for foreign key and CHECK constraints. If the constraint is marked NOT VALID, the potentially-lengthy initial check to verify that all rows in the table satisfy the constraint is skipped. The constraint will still be enforced against subsequent inserts or updates (that is, they'll fail unless there is a matching row in the referenced table, in the case of foreign keys; and they'll fail unless the new row matches the specified check constraints). But the database will not assume that the constraint holds for all rows in the table, until it is validated by using the VALIDATE CONSTRAINT option. Foreign key constraints on partitioned tables may not be declared NOT VALID at present.

The addition of a foreign key constraint requires a SHARE ROW EXCLUSIVE lock on the referenced table.

Additional restrictions apply when unique or primary key constraints are added to partitioned tables; see CREATE TABLE.

ADD table_constraint_using_index

This form adds a new PRIMARY KEY or UNIQUE constraint to a table based on an existing unique index. All the columns of the index will be included in the constraint.

The index cannot have expression columns nor be a partial index. Also, it must be a b-tree index with default sort ordering. These restrictions ensure that the index is equivalent to one that would be built by a regular ADD PRIMARY KEY or ADD UNIQUE command.

If PRIMARY KEY is specified, and the index's columns are not already marked NOT NULL, then this command will attempt to do ALTER COLUMN SET NOT NULL against each such column. That requires a full table scan to verify the column(s) contain no nulls. In all other cases, this is a fast operation.

If a constraint name is provided then the index will be renamed to match the constraint name. Otherwise the constraint will be named the same as the index.

After this command is executed, the index is “owned” by the constraint, in the same way as if the index had been built by a regular ADD PRIMARY KEY or ADD UNIQUE command. In particular, dropping the constraint will make the index disappear too.

This form is not currently supported on partitioned tables.

Note

Adding a constraint using an existing index can be helpful in situations where a new constraint needs to be added without blocking table updates for a long time.
ALTER TABLE

To do that, create the index using CREATE INDEX CONCURRENTLY, and then install it as an official constraint using this syntax. See the example below.

ALTER CONSTRAINT

This form alters the attributes of a constraint that was previously created. Currently only foreign key constraints may be altered.

VALIDATE CONSTRAINT

This form validates a foreign key or check constraint that was previously created as NOT VALID, by scanning the table to ensure there are no rows for which the constraint is not satisfied. Nothing happens if the constraint is already marked valid.

Validation can be a long process on larger tables. The value of separating validation from initial creation is that you can defer validation to less busy times, or can be used to give additional time to correct pre-existing errors while preventing new errors. Note also that validation on its own does not prevent normal write commands against the table while it runs.

Validation acquires only a SHARE UPDATE EXCLUSIVE lock on the table being altered. If the constraint is a foreign key then a ROW SHARE lock is also required on the table referenced by the constraint.

DROP CONSTRAINT [ IF EXISTS ]

This form drops the specified constraint on a table, along with any index underlying the constraint. If IF EXISTS is specified and the constraint does not exist, no error is thrown. In this case a notice is issued instead.

DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER

These forms configure the firing of trigger(s) belonging to the table. A disabled trigger is still known to the system, but is not executed when its triggering event occurs. For a deferred trigger, the enable status is checked when the event occurs, not when the trigger function is actually executed. One can disable or enable a single trigger specified by name, or all triggers on the table, or only user triggers (this option excludes internally generated constraint triggers such as those that are used to implement foreign key constraints or deferrable uniqueness and exclusion constraints). Disabling or enabling internally generated constraint triggers requires superuser privileges; it should be done with caution since of course the integrity of the constraint cannot be guaranteed if the triggers are not executed.

The trigger firing mechanism is also affected by the configuration variable session_replication_role. Simply enabled triggers (the default) will fire when the replication role is “origin” (the default) or “local”. Triggers configured as ENABLE REPLICA will only fire if the session is in “replica” mode, and triggers configured as ENABLE ALWAYS will fire regardless of the current replication role.

The effect of this mechanism is that in the default configuration, triggers do not fire on replicas. This is useful because if a trigger is used on the origin to propagate data between tables, then the replication system will also replicate the propagated data, and the trigger should not fire a second time on the replica, because that would lead to duplication. However, if a trigger is used for another purpose such as creating external alerts, then it might be appropriate to set it to ENABLE ALWAYS so that it is also fired on replicas.

This command acquires a SHARE ROW EXCLUSIVE lock.

DISABLE/ENABLE [ REPLICA | ALWAYS ] RULE

These forms configure the firing of rewrite rules belonging to the table. A disabled rule is still known to the system, but is not applied during query rewriting. The semantics are as for dis-
ALTER TABLE

abled/enabled triggers. This configuration is ignored for ON SELECT rules, which are always applied in order to keep views working even if the current session is in a non-default replication role.

The rule firing mechanism is also affected by the configuration variable session_replication_role, analogous to triggers as described above.

DISABLE/ENABLE ROW LEVEL SECURITY

These forms control the application of row security policies belonging to the table. If enabled and no policies exist for the table, then a default-deny policy is applied. Note that policies can exist for a table even if row level security is disabled - in this case, the policies will NOT be applied and the policies will be ignored. See also CREATE POLICY.

NO FORCE/FORCE ROW LEVEL SECURITY

These forms control the application of row security policies belonging to the table when the user is the table owner. If enabled, row level security policies will be applied when the user is the table owner. If disabled (the default) then row level security will not be applied when the user is the table owner. See also CREATE POLICY.

CLUSTER ON

This form selects the default index for future CLUSTER operations. It does not actually re-cluster the table.

Changing cluster options acquires a SHARE UPDATE EXCLUSIVE lock.

SET WITHOUT CLUSTER

This form removes the most recently used CLUSTER index specification from the table. This affects future cluster operations that don't specify an index.

Changing cluster options acquires a SHARE UPDATE EXCLUSIVE lock.

SET WITH OIDS

This form adds an oid system column to the table (see Section 5.4). It does nothing if the table already has OIDs.

Note that this is not equivalent to ADD COLUMN oid oid; that would add a normal column that happened to be named oid, not a system column.

SET WITHOUT OIDS

This form removes the oid system column from the table. This is exactly equivalent to DROP COLUMN oid RESTRICT, except that it will not complain if there is already no oid column.

SET TABLESPACE

This form changes the table's tablespace to the specified tablespace and moves the data file(s) associated with the table to the new tablespace. Indexes on the table, if any, are not moved; but they can be moved separately with additional SET TABLESPACE commands. All tables in the current database in a tablespace can be moved by using the ALL IN TABLESPACE form, which will lock all tables to be moved first and then move each one. This form also supports OWNED BY, which will only move tables owned by the roles specified. If the NOWAIT option is specified then the command will fail if it is unable to acquire all of the locks required immediately. Note that system catalogs are not moved by this command, use ALTER DATABASE or explicit ALTER TABLE invocations instead if desired. The information_schema relations are not considered part of the system catalogs and will be moved. See also CREATE TABLESPACE.
ALTER TABLE

SET { LOGGED | UNLOGGED }

This form changes the table from unlogged to logged or vice-versa (see UNLOGGED). It cannot be applied to a temporary table.

SET ( storage_parameter = value [, ... ] )

This form changes one or more storage parameters for the table. See Storage Parameters for details on the available parameters. Note that the table contents will not be modified immediately by this command; depending on the parameter you might need to rewrite the table to get the desired effects. That can be done with VACUUM FULL, CLUSTER or one of the forms of ALTER TABLE that forces a table rewrite. For planner related parameters, changes will take effect from the next time the table is locked so currently executing queries will not be affected.

SHARE UPDATE EXCLUSIVE lock will be taken for fillfactor, toast and autovacuum storage parameters, as well as the following planner related parameters: effective_io_concurrency, parallel_workers, seq_page_cost, random_page_cost, n_distinct and n_distinct_inherited.

Note

While CREATE TABLE allows OIDS to be specified in the WITH (storage_parameter) syntax, ALTER TABLE does not treat OIDS as a storage parameter. Instead use the SET WITH OIDS and SET WITHOUT OIDS forms to change OID status.

RESET ( storage_parameter [, ... ] )

This form resets one or more storage parameters to their defaults. As with SET, a table rewrite might be needed to update the table entirely.

INHERIT parent_table

This form adds the target table as a new child of the specified parent table. Subsequently, queries against the parent will include records of the target table. To be added as a child, the target table must already contain all the same columns as the parent (it could have additional columns, too). The columns must have matching data types, and if they have NOT NULL constraints in the parent then they must also have NOT NULL constraints in the child.

There must also be matching child-table constraints for all CHECK constraints of the parent, except those marked non-inheritable (that is, created with ALTER TABLE ... ADD CONSTRAINT ... NO INHERIT) in the parent, which are ignored; all child-table constraints matched must not be marked non-inheritable. Currently UNIQUE, PRIMARY KEY, and FOREIGN KEY constraints are not considered, but this might change in the future.

NO INHERIT parent_table

This form removes the target table from the list of children of the specified parent table. Queries against the parent table will no longer include records drawn from the target table.

OF type_name

This form links the table to a composite type as though CREATE TABLE OF had formed it. The table's list of column names and types must precisely match that of the composite type; the presence of an oid system column is permitted to differ. The table must not inherit from any other table. These restrictions ensure that CREATE TABLE OF would permit an equivalent table definition.
ALTER TABLE

NOT OF

This form dissociates a typed table from its type.

OWNER TO

This form changes the owner of the table, sequence, view, materialized view, or foreign table to the specified user.

REPLICA IDENTITY

This form changes the information which is written to the write-ahead log to identify rows which are updated or deleted. This option has no effect except when logical replication is in use. DEFAULT (the default for non-system tables) records the old values of the columns of the primary key, if any. USING INDEX records the old values of the columns covered by the named index, which must be unique, not partial, not deferrable, and include only columns marked NOT NULL. FULL records the old values of all columns in the row. NOTHING records no information about the old row. (This is the default for system tables.) In all cases, no old values are logged unless at least one of the columns that would be logged differs between the old and new versions of the row.

RENAME

The RENAME forms change the name of a table (or an index, sequence, view, materialized view, or foreign table), the name of an individual column in a table, or the name of a constraint of the table. When renaming a constraint that has an underlying index, the index is renamed as well. There is no effect on the stored data.

SET SCHEMA

This form moves the table into another schema. Associated indexes, constraints, and sequences owned by table columns are moved as well.

ATTACH PARTITION partition_name { FOR VALUES partition_bound_spec | DEFAULT }

This form attaches an existing table (which might itself be partitioned) as a partition of the target table. The table can be attached as a partition for specific values using FOR VALUES or as a default partition by using DEFAULT. For each index in the target table, a corresponding one will be created in the attached table; or, if an equivalent index already exists, will be attached to the target table’s index, as if ALTER INDEX ATTACH PARTITION had been executed.

A partition using FOR VALUES uses same syntax for partition_bound_spec as CREATE TABLE. The partition bound specification must correspond to the partitioning strategy and partition key of the target table. The table to be attached must have all the same columns as the target table and no more; moreover, the column types must also match. Also, it must have all the NOT NULL and CHECK constraints of the target table. Currently FOREIGN KEY constraints are not considered. UNIQUE and PRIMARY KEY constraints from the parent table will be created in the partition, if they don’t already exist. If any of the CHECK constraints of the table being attached is marked NO INHERIT, the command will fail; such a constraint must be recreated without the NO INHERIT clause.

If the new partition is a regular table, a full table scan is performed to check that no existing row in the table violates the partition constraint. It is possible to avoid this scan by adding a valid CHECK constraint to the table that would allow only the rows satisfying the desired partition constraint before running this command. It will be determined using such a constraint that the table need not be scanned to validate the partition constraint. This does not work, however, if any of the partition keys is an expression and the partition does not accept NULL values. If attaching a list partition that will not accept NULL values, also add NOT NULL constraint to the partition key column, unless it’s an expression.
If the new partition is a foreign table, nothing is done to verify that all the rows in the foreign table obey the partition constraint. (See the discussion in CREATE FOREIGN TABLE about constraints on the foreign table.)

When a table has a default partition, defining a new partition changes the partition constraint for the default partition. The default partition can’t contain any rows that would need to be moved to the new partition, and will be scanned to verify that none are present. This scan, like the scan of the new partition, can be avoided if an appropriate CHECK constraint is present. Also like the scan of the new partition, it is always skipped when the default partition is a foreign table.

DETACH PARTITION partition_name

This form detaches specified partition of the target table. The detached partition continues to exist as a standalone table, but no longer has any ties to the table from which it was detached. Any indexes that were attached to the target table's indexes are detached.

All the forms of ALTER TABLE that act on a single table, except RENAME, SET SCHEMA, ATTACH PARTITION, and DETACH PARTITION can be combined into a list of multiple alterations to be applied together. For example, it is possible to add several columns and/or alter the type of several columns in a single command. This is particularly useful with large tables, since only one pass over the table need be made.

You must own the table to use ALTER TABLE. To change the schema or tablespace of a table, you must also have CREATE privilege on the new schema or tablespace. To add the table as a new child of a parent table, you must own the parent table as well. Also, to attach a table as a new partition of the table, you must own the table being attached. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the table's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the table. However, a superuser can alter ownership of any table anyway.) To add a column or alter a column type or use the OF clause, you must also have USAGE privilege on the data type.

**Parameters**

**IF EXISTS**

Do not throw an error if the table does not exist. A notice is issued in this case.

*name*

The name (optionally schema-qualified) of an existing table to alter. If ONLY is specified before the table name, only that table is altered. If ONLY is not specified, the table and all its descendant tables (if any) are altered. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

*column_name*

Name of a new or existing column.

*new_column_name*

New name for an existing column.

*new_name*

New name for the table.

*data_type*

Data type of the new column, or new data type for an existing column.
ALTER TABLE

table_constraint

    New table constraint for the table.

constraint_name

    Name of a new or existing constraint.

CASCADE

    Automatically drop objects that depend on the dropped column or constraint (for example, views referencing the column), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

    Refuse to drop the column or constraint if there are any dependent objects. This is the default behavior.

trigger_name

    Name of a single trigger to disable or enable.

ALL

    Disable or enable all triggers belonging to the table. (This requires superuser privilege if any of the triggers are internally generated constraint triggers such as those that are used to implement foreign key constraints or deferrable uniqueness and exclusion constraints.)

USER

    Disable or enable all triggers belonging to the table except for internally generated constraint triggers such as those that are used to implement foreign key constraints or deferrable uniqueness and exclusion constraints.

index_name

    The name of an existing index.

storage_parameter

    The name of a table storage parameter.

value

    The new value for a table storage parameter. This might be a number or a word depending on the parameter.

parent_table

    A parent table to associate or de-associate with this table.

new_owner

    The user name of the new owner of the table.

new_tablespace

    The name of the tablespace to which the table will be moved.

new_schema

    The name of the schema to which the table will be moved.
ALTER TABLE

`partition_name`

The name of the table to attach as a new partition or to detach from this table.

`partition_bound_spec`

The partition bound specification for a new partition. Refer to CREATE TABLE for more details on the syntax of the same.

Notes

The key word COLUMN is noise and can be omitted.

When a column is added with ADD COLUMN and a non-volatile DEFAULT is specified, the default is evaluated at the time of the statement and the result stored in the table's metadata. That value will be used for the column for all existing rows. If no DEFAULT is specified, NULL is used. In neither case is a rewrite of the table required.

Adding a column with a volatile DEFAULT or changing the type of an existing column will require the entire table and its indexes to be rewritten. As an exception, when changing the type of an existing column, if the USING clause does not change the column contents and the old type is either binary coercible to the new type or an unconstrained domain over the new type, a table rewrite is not needed; but any indexes on the affected columns must still be rebuilt. Adding or removing a system oid column also requires rewriting the entire table. Table and/or index rebuilds may take a significant amount of time for a large table; and will temporarily require as much as double the disk space.

Adding a CHECK or NOT NULL constraint requires scanning the table to verify that existing rows meet the constraint, but does not require a table rewrite.

Similarly, when attaching a new partition it may be scanned to verify that existing rows meet the partition constraint.

The main reason for providing the option to specify multiple changes in a single ALTER TABLE is that multiple table scans or rewrites can thereby be combined into a single pass over the table.

The DROP COLUMN form does not physically remove the column, but simply makes it invisible to SQL operations. Subsequent insert and update operations in the table will store a null value for the column. Thus, dropping a column is quick but it will not immediately reduce the on-disk size of your table, as the space occupied by the dropped column is not reclaimed. The space will be reclaimed over time as existing rows are updated. (These statements do not apply when dropping the system oid column; that is done with an immediate rewrite.)

To force immediate reclamation of space occupied by a dropped column, you can execute one of the forms of ALTER TABLE that performs a rewrite of the whole table. This results in reconstructing each row with the dropped column replaced by a null value.

The rewriting forms of ALTER TABLE are not MVCC-safe. After a table rewrite, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the rewrite occurred. See Section 13.5 for more details.

The USING option of SET DATA TYPE can actually specify any expression involving the old values of the row; that is, it can refer to other columns as well as the one being converted. This allows very general conversions to be done with the SET DATA TYPE syntax. Because of this flexibility, the USING expression is not applied to the column’s default value (if any); the result might not be a constant expression as required for a default. This means that when there is no implicit or assignment cast from old to new type, SET DATA TYPE might fail to convert the default even though a USING clause is supplied. In such cases, drop the default with DROP DEFAULT, perform the ALTER TYPE, and then use SET DEFAULT to add a suitable new default. Similar considerations apply to indexes and constraints involving the column.
If a table has any descendant tables, it is not permitted to add, rename, or change the type of a column in the parent table without doing the same to the descendants. This ensures that the descendants always have columns matching the parent. Similarly, a CHECK constraint cannot be renamed in the parent without also renaming it in all descendants, so that CHECK constraints also match between the parent and its descendants. (That restriction does not apply to index-based constraints, however.) Also, because selecting from the parent also selects from its descendants, a constraint on the parent cannot be marked valid unless it is also marked valid for those descendants. In all of these cases, ALTER TABLE ONLY will be rejected.

A recursive DROP COLUMN operation will remove a descendant table’s column only if the descendant does not inherit that column from any other parents and never had an independent definition of the column. A nonrecursive DROP COLUMN (i.e., ALTER TABLE ONLY ... DROP COLUMN) never removes any descendant columns, but instead marks them as independently defined rather than inherited. A nonrecursive DROP COLUMN command will fail for a partitioned table, because all partitions of a table must have the same columns as the partitioning root.

The actions for identity columns (ADD GENERATED, SET etc., DROP IDENTITY), as well as the actions TRIGGER, CLUSTER, OWNER, and TABLESPACE never recurse to descendant tables; that is, they always act as though ONLY were specified. Adding a constraint recurses only for CHECK constraints that are not marked NO INHERIT.

Changing any part of a system catalog table is not permitted.

Refer to CREATE TABLE for a further description of valid parameters. Chapter 5 has further information on inheritance.

**Examples**

To add a column of type varchar to a table:

```
ALTER TABLE distributors ADD COLUMN address varchar(30);
```

To drop a column from a table:

```
ALTER TABLE distributors DROP COLUMN address RESTRICT;
```

To change the types of two existing columns in one operation:

```
ALTER TABLE distributors
    ALTER COLUMN address TYPE varchar(80),
    ALTER COLUMN name TYPE varchar(100);
```

To change an integer column containing Unix timestamps to timestamp with time zone via a USING clause:

```
ALTER TABLE foo
    ALTER COLUMN foo_timestamp SET DATA TYPE timestamp with time zone
    USING
        timestamp with time zone 'epoch' + foo_timestamp * interval '1 second';
```

The same, when the column has a default expression that won’t automatically cast to the new data type:

```
ALTER TABLE foo
```
ALTER TABLE

ALTER COLUMN foo_timestamp DROP DEFAULT,
ALTER COLUMN foo_timestamp TYPE timestamp with time zone USING
timestamp with time zone 'epoch' + foo_timestamp * interval '1 second',
ALTER COLUMN foo_timestamp SET DEFAULT now();

To rename an existing column:

ALTER TABLE distributors RENAME COLUMN address TO city;

To rename an existing table:

ALTER TABLE distributors RENAME TO suppliers;

To rename an existing constraint:

ALTER TABLE distributors RENAME CONSTRAINT zipchk TO zip_check;

To add a not-null constraint to a column:

ALTER TABLE distributors ALTER COLUMN street SET NOT NULL;

To remove a not-null constraint from a column:

ALTER TABLE distributors ALTER COLUMN street DROP NOT NULL;

To add a check constraint to a table and all its children:

ALTER TABLE distributors ADD CONSTRAINT zipchk CHECK
(char_length(zipcode) = 5);

To add a check constraint only to a table and not to its children:

ALTER TABLE distributors ADD CONSTRAINT zipchk CHECK
(char_length(zipcode) = 5) NO INHERIT;

(The check constraint will not be inherited by future children, either.)

To remove a check constraint from a table and all its children:

ALTER TABLE distributors DROP CONSTRAINT zipchk;

To remove a check constraint from one table only:

ALTER TABLE ONLY distributors DROP CONSTRAINT zipchk;

(The check constraint remains in place for any child tables.)

To add a foreign key constraint to a table:
ALTER TABLE distributors ADD CONSTRAINT distfk FOREIGN KEY (address) REFERENCES addresses (address);

To add a foreign key constraint to a table with the least impact on other work:

ALTER TABLE distributors ADD CONSTRAINT distfk FOREIGN KEY (address) REFERENCES addresses (address) NOT VALID;
ALTER TABLE distributors VALIDATE CONSTRAINT distfk;

To add a (multicolumn) unique constraint to a table:

ALTER TABLE distributors ADD CONSTRAINT dist_id_zipcode_key UNIQUE (dist_id, zipcode);

To add an automatically named primary key constraint to a table, noting that a table can only ever have one primary key:

ALTER TABLE distributors ADD PRIMARY KEY (dist_id);

To move a table to a different tablespace:

ALTER TABLE distributors SET TABLESPACE fasttablespace;

To move a table to a different schema:

ALTER TABLE myschema.distributors SET SCHEMA yourschema;

To recreate a primary key constraint, without blocking updates while the index is rebuilt:

CREATE UNIQUE INDEX CONCURRENTLY dist_id_temp_idx ON distributors (dist_id);
ALTER TABLE distributors DROP CONSTRAINT distributors_pkey, ADD CONSTRAINT distributors_pkey PRIMARY KEY USING INDEX dist_id_temp_idx;

To attach a partition to a range-partitioned table:

ALTER TABLE measurement
    ATTACH PARTITION measurement_y2016m07 FOR VALUES FROM ('2016-07-01') TO ('2016-08-01');

To attach a partition to a list-partitioned table:

ALTER TABLE cities
    ATTACH PARTITION cities_ab FOR VALUES IN ('a', 'b');

To attach a partition to a hash-partitioned table:

ALTER TABLE orders
    ATTACH PARTITION orders_p4 FOR VALUES WITH (MODULUS 4, REMAINDER 3);

To attach a default partition to a partitioned table:
ALTER TABLE cities
   ATTACH PARTITION cities_partdef DEFAULT;

To detach a partition from a partitioned table:

ALTER TABLE measurement
   DETACH PARTITION measurement_y2015m12;

Compatibility

The forms ADD (without USING INDEX), DROP [COLUMN], DROP IDENTITY, RESTART, SET DEFAULT, SET DATA TYPE (without USING), SET GENERATED, and SET sequence_option conform with the SQL standard. The other forms are PostgreSQL extensions of the SQL standard. Also, the ability to specify more than one manipulation in a single ALTER TABLE command is an extension.

ALTER TABLE DROP COLUMN can be used to drop the only column of a table, leaving a zero-column table. This is an extension of SQL, which disallows zero-column tables.

See Also

CREATE TABLE
ALTER TABLESPACE

ALTER TABLESPACE — change the definition of a tablespace

Synopsis

ALTER TABLESPACE name RENAME TO new_name
ALTER TABLESPACE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER TABLESPACE name SET ( tablespace_option = value [, ... ] )
ALTER TABLESPACE name RESET ( tablespace_option [, ... ] )

Description

ALTER TABLESPACE can be used to change the definition of a tablespace.

You must own the tablespace to change the definition of a tablespace. To alter the owner, you must also be a direct or indirect member of the new owning role. (Note that superusers have these privileges automatically.)

Parameters

name

The name of an existing tablespace.

new_name

The new name of the tablespace. The new name cannot begin with pg_, as such names are reserved for system tablespaces.

new_owner

The new owner of the tablespace.

tablespace_option

A tablespace parameter to be set or reset. Currently, the only available parameters are seq_page_cost, random_page_cost and effective_io_concurrency. Setting either value for a particular tablespace will override the planner's usual estimate of the cost of reading pages from tables in that tablespace, as established by the configuration parameters of the same name (see seq_page_cost, random_page_cost, effective_io_concurrency). This may be useful if one tablespace is located on a disk which is faster or slower than the remainder of the I/O sub-system.

Examples

Rename tablespace index_space to fast_raid:

ALTER TABLESPACE index_space RENAME TO fast_raid;

Change the owner of tablespace index_space:

ALTER TABLESPACE index_space OWNER TO new_owner;
ALTER TABLESPACE

ALTER TABLESPACE index_space OWNER TO mary;

Compatibility

There is no ALTER TABLESPACE statement in the SQL standard.

See Also

CREATE TABLESPACE, DROP TABLESPACE
ALTER TEXT SEARCH CONFIGURATION

ALTER TEXT SEARCH CONFIGURATION — change the definition of a text search configuration

Synopsis

ALTER TEXT SEARCH CONFIGURATION name
       ADD MAPPING FOR token_type [, ... ] WITH dictionary_name [, ... ]
ALTER TEXT SEARCH CONFIGURATION name
       ALTER MAPPING FOR token_type [, ... ] WITH dictionary_name [, ... ]
ALTER TEXT SEARCH CONFIGURATION name
       ALTER MAPPING REPLACE old_dictionary WITH new_dictionary
ALTER TEXT SEARCH CONFIGURATION name
       ALTER MAPPING FOR token_type [, ... ] REPLACE old_dictionary WITH new_dictionary
ALTER TEXT SEARCH CONFIGURATION name
       DROP MAPPING [ IF EXISTS ] FOR token_type [, ... ]
ALTER TEXT SEARCH CONFIGURATION name RENAME TO new_name
ALTER TEXT SEARCH CONFIGURATION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER TEXT SEARCH CONFIGURATION name SET SCHEMA new_schema

Description

ALTER TEXT SEARCH CONFIGURATION changes the definition of a text search configuration.
You can modify its mappings from token types to dictionaries, or change the configuration's name
or owner.

You must be the owner of the configuration to use ALTER TEXT SEARCH CONFIGURATION.

Parameters

name

The name (optionally schema-qualified) of an existing text search configuration.

token_type

The name of a token type that is emitted by the configuration's parser.

dictionary_name

The name of a text search dictionary to be consulted for the specified token type(s). If multiple
dictionaries are listed, they are consulted in the specified order.

old_dictionary

The name of a text search dictionary to be replaced in the mapping.

new_dictionary

The name of a text search dictionary to be substituted for old_dictionary.
new_name

The new name of the text search configuration.

new_owner

The new owner of the text search configuration.

new_schema

The new schema for the text search configuration.

The **ADD MAPPING FOR** form installs a list of dictionaries to be consulted for the specified token type(s); it is an error if there is already a mapping for any of the token types. The **ALTER MAPPING FOR** form does the same, but first removing any existing mapping for those token types. The **ALTER MAPPING REPLACE** forms substitute `new_dictionary` for `old_dictionary` anywhere the latter appears. This is done for only the specified token types when `FOR` appears, or for all mappings of the configuration when it doesn't. The **DROP MAPPING** form removes all dictionaries for the specified token type(s), causing tokens of those types to be ignored by the text search configuration. It is an error if there is no mapping for the token types, unless **IF EXISTS** appears.

**Examples**

The following example replaces the `english` dictionary with the `swedish` dictionary anywhere that `english` is used within `my_config`.

```
ALTER TEXT SEARCH CONFIGURATION my_config
  ALTER MAPPING REPLACE english WITH swedish;
```

**Compatibility**

There is no **ALTER TEXT SEARCH CONFIGURATION** statement in the SQL standard.

**See Also**

CREATE TEXT SEARCH CONFIGURATION, DROP TEXT SEARCH CONFIGURATION
ALTER TEXT SEARCH DICTIONARY

ALTER TEXT SEARCH DICTIONARY — change the definition of a text search dictionary

Synopsis

```
ALTER TEXT SEARCH DICTIONARY name (  
   option [ = value ] [, ... ]  
)  
ALTER TEXT SEARCH DICTIONARY name RENAME TO new_name  
ALTER TEXT SEARCH DICTIONARY name OWNER TO { new_owner |  
   CURRENT_USER | SESSION_USER }  
ALTER TEXT SEARCH DICTIONARY name SET SCHEMA new_schema
```

Description

ALTER TEXT SEARCH DICTIONARY changes the definition of a text search dictionary. You can change the dictionary's template-specific options, or change the dictionary's name or owner.

You must be the owner of the dictionary to use ALTER TEXT SEARCH DICTIONARY.

Parameters

name

The name (optionally schema-qualified) of an existing text search dictionary.

option

The name of a template-specific option to be set for this dictionary.

value

The new value to use for a template-specific option. If the equal sign and value are omitted, then any previous setting for the option is removed from the dictionary, allowing the default to be used.

new_name

The new name of the text search dictionary.

new_owner

The new owner of the text search dictionary.

new_schema

The new schema for the text search dictionary.

Template-specific options can appear in any order.

Examples

The following example command changes the stopword list for a Snowball-based dictionary. Other parameters remain unchanged.
ALTER TEXT SEARCH DICTIONARY my_dict ( StopWords = newrussian );

The following example command changes the language option to dutch, and removes the stopword option entirely.

ALTER TEXT SEARCH DICTIONARY my_dict ( language = dutch, StopWords );

The following example command “updates” the dictionary’s definition without actually changing anything.

ALTER TEXT SEARCH DICTIONARY my_dict ( dummy );

(The reason this works is that the option removal code doesn’t complain if there is no such option.) This trick is useful when changing configuration files for the dictionary: the ALTER will force existing database sessions to re-read the configuration files, which otherwise they would never do if they had read them earlier.

Compatibility

There is no ALTER TEXT SEARCH DICTIONARY statement in the SQL standard.

See Also

CREATE TEXT SEARCH DICTIONARY, DROP TEXT SEARCH DICTIONARY
ALTER TEXT SEARCH PARSER

ALTER TEXT SEARCH PARSER — change the definition of a text search parser

Synopsis

```
ALTER TEXT SEARCH PARSER name RENAME TO new_name
ALTER TEXT SEARCH PARSER name SET SCHEMA new_schema
```

Description

ALTER TEXT SEARCH PARSER changes the definition of a text search parser. Currently, the only supported functionality is to change the parser's name.

You must be a superuser to use ALTER TEXT SEARCH PARSER.

Parameters

`name`

The name (optionally schema-qualified) of an existing text search parser.

`new_name`

The new name of the text search parser.

`new_schema`

The new schema for the text search parser.

Compatibility

There is no ALTER TEXT SEARCH PARSER statement in the SQL standard.

See Also

CREATE TEXT SEARCH PARSER, DROP TEXT SEARCH PARSER
ALTER TEXT SEARCH TEMPLATE

ALTER TEXT SEARCH TEMPLATE — change the definition of a text search template

Synopsis

ALTER TEXT SEARCH TEMPLATE name RENAME TO new_name
ALTER TEXT SEARCH TEMPLATE name SET SCHEMA new_schema

Description

ALTER TEXT SEARCH TEMPLATE changes the definition of a text search template. Currently, the only supported functionality is to change the template's name.

You must be a superuser to use ALTER TEXT SEARCH TEMPLATE.

Parameters

name

The name (optionally schema-qualified) of an existing text search template.

new_name

The new name of the text search template.

new_schema

The new schema for the text search template.

Compatibility

There is no ALTER TEXT SEARCH TEMPLATE statement in the SQL standard.

See Also

CREATE TEXT SEARCH TEMPLATE, DROP TEXT SEARCH TEMPLATE
ALTER TRIGGER

ALTER TRIGGER — change the definition of a trigger

Synopsis

ALTER TRIGGER name ON table_name RENAME TO new_name
ALTER TRIGGER name ON table_name DEPENDS ON
   EXTENSION extension_name

Description

ALTER TRIGGER changes properties of an existing trigger. The RENAME clause changes the name of the given trigger without otherwise changing the trigger definition. The DEPENDS ON EXTENSION clause marks the trigger as dependent on an extension, such that if the extension is dropped, the trigger will automatically be dropped as well.

You must own the table on which the trigger acts to be allowed to change its properties.

Parameters

name
   The name of an existing trigger to alter.

table_name
   The name of the table on which this trigger acts.

new_name
   The new name for the trigger.

extension_name
   The name of the extension that the trigger is to depend on.

Notes

The ability to temporarily enable or disable a trigger is provided by ALTER TABLE, not by ALTER TRIGGER, because ALTER TRIGGER has no convenient way to express the option of enabling or disabling all of a table's triggers at once.

Examples

To rename an existing trigger:

   ALTER TRIGGER emp_stamp ON emp RENAME TO emp_track_chgs;

To mark a trigger as being dependent on an extension:

   ALTER TRIGGER emp_stamp ON emp DEPENDS ON EXTENSION emplib;
Compatibility

ALTER TRIGGER is a PostgreSQL extension of the SQL standard.

See Also

ALTER TABLE
ALTER TYPE

ALTER TYPE — change the definition of a type

Synopsis

ALTER TYPE name action [, ... ]
ALTER TYPE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER TYPE name RENAME ATTRIBUTE attribute_name
   TO new_attribute_name [ CASCADE | RESTRICT ]
ALTER TYPE name RENAME TO new_name
ALTER TYPE name SET SCHEMA new_schema
ALTER TYPE name ADD VALUE [ IF NOT EXISTS ] new_enum_value
   [ { BEFORE | AFTER } neighbor_enum_value ]
ALTER TYPE name RENAME VALUE existing_enum_value TO new_enum_value

where action is one of:

   ADD ATTRIBUTE attribute_name data_type [ COLLATE collation ]
   [ CASCADE | RESTRICT ]
   DROP ATTRIBUTE [ IF EXISTS ] attribute_name [ CASCADE | RESTRICT ]
   ALTER ATTRIBUTE attribute_name [ SET DATA ] TYPE data_type
   [ COLLATE collation ] [ CASCADE | RESTRICT ]

Description

ALTER TYPE changes the definition of an existing type. There are several subforms:

ADD ATTRIBUTE

This form adds a new attribute to a composite type, using the same syntax as CREATE TYPE.

DROP ATTRIBUTE [ IF EXISTS ]

This form drops an attribute from a composite type. If IF EXISTS is specified and the attribute does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of an attribute of a composite type.

OWNER

This form changes the owner of the type.

RENAME

This form changes the name of the type or the name of an individual attribute of a composite type.

SET SCHEMA

This form moves the type into another schema.
ALTER TYPE

ADD VALUE [ IF NOT EXISTS ] [ BEFORE | AFTER ]

This form adds a new value to an enum type. The new value's place in the enum's ordering can be specified as being BEFORE or AFTER one of the existing values. Otherwise, the new item is added at the end of the list of values.

If IF NOT EXISTS is specified, it is not an error if the type already contains the new value: a notice is issued but no other action is taken. Otherwise, an error will occur if the new value is already present.

RENAME VALUE

This form renames a value of an enum type. The value's place in the enum's ordering is not affected. An error will occur if the specified value is not present or the new name is already present.

The ADD ATTRIBUTE, DROP ATTRIBUTE, and ALTER ATTRIBUTE actions can be combined into a list of multiple alterations to apply in parallel. For example, it is possible to add several attributes and/or alter the type of several attributes in a single command.

You must own the type to use ALTER TYPE. To change the schema of a type, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the type's schema. (These restrictions enforce that altering the owner doesn't do anything you couldn't do by dropping and recreating the type. However, a superuser can alter ownership of any type anyway.) To add an attribute or alter an attribute type, you must also have USAGE privilege on the data type.

Parameters

name

The name (possibly schema-qualified) of an existing type to alter.

new_name

The new name for the type.

new_owner

The user name of the new owner of the type.

new_schema

The new schema for the type.

attribute_name

The name of the attribute to add, alter, or drop.

new_attribute_name

The new name of the attribute to be renamed.

data_type

The data type of the attribute to add, or the new type of the attribute to alter.

new_enum_value

The new value to be added to an enum type's list of values, or the new name to be given to an existing value. Like all enum literals, it needs to be quoted.
ALTER TYPE

`neighbor_enum_value`

The existing enum value that the new value should be added immediately before or after in the enum type's sort ordering. Like all enum literals, it needs to be quoted.

`existing_enum_value`

The existing enum value that should be renamed. Like all enum literals, it needs to be quoted.

`CASCADE`

Automatically propagate the operation to typed tables of the type being altered, and their descendants.

`RESTRICT`

Refuse the operation if the type being altered is the type of a typed table. This is the default.

**Notes**

ALTER TYPE ... ADD VALUE (the form that adds a new value to an enum type) cannot be executed inside a transaction block.

Comparisons involving an added enum value will sometimes be slower than comparisons involving only original members of the enum type. This will usually only occur if BEFORE or AFTER is used to set the new value's sort position somewhere other than at the end of the list. However, sometimes it will happen even though the new value is added at the end (this occurs if the OID counter "wrapped around" since the original creation of the enum type). The slowdown is usually insignificant; but if it matters, optimal performance can be regained by dropping and recreating the enum type, or by dumping and reloading the database.

**Examples**

To rename a data type:

```
ALTER TYPE electronic_mail RENAME TO email;
```

To change the owner of the type email to joe:

```
ALTER TYPE email OWNER TO joe;
```

To change the schema of the type email to customers:

```
ALTER TYPE email SET SCHEMA customers;
```

To add a new attribute to a type:

```
ALTER TYPE compfoo ADD ATTRIBUTE f3 int;
```

To add a new value to an enum type in a particular sort position:

```
ALTER TYPE colors ADD VALUE 'orange' AFTER 'red';
```

To rename an enum value:
ALTER TYPE colors RENAME VALUE 'purple' TO 'mauve';

**Compatibility**

The variants to add and drop attributes are part of the SQL standard; the other variants are PostgreSQL extensions.

**See Also**

CREATE TYPE, DROP TYPE
ALTER USER

ALTER USER — change a database role

Synopsis

ALTER USER role_specification [ WITH ] option [ ... ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICAUTION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT connlimit
- [ ENCRYPTED ] PASSWORD 'password'
- VALID UNTIL 'timestamp'

ALTER USER name RENAME TO new_name

ALTER USER { role_specification | ALL } [ IN DATABASE database_name ]
SET configuration_parameter { TO | = } { value | DEFAULT }
ALTER USER { role_specification | ALL } [ IN DATABASE database_name ]
SET configuration_parameter FROM CURRENT
ALTER USER { role_specification | ALL } [ IN DATABASE database_name ]
RESET configuration_parameter
ALTER USER { role_specification | ALL } [ IN DATABASE database_name ]
RESET ALL

where role_specification can be:

- role_name
- CURRENT_USER
- SESSION_USER

Description

ALTER USER is now an alias for ALTER ROLE.

Compatibility

The ALTER USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also

ALTER ROLE
ALTER USER MAPPING

ALTER USER MAPPING — change the definition of a user mapping

Synopsis

ALTER USER MAPPING FOR { user_name | USER | CURRENT_USER | SESSION_USER | PUBLIC } SERVER server_name OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )

Description

ALTER USER MAPPING changes the definition of a user mapping.

The owner of a foreign server can alter user mappings for that server for any user. Also, a user can alter a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters

user_name

User name of the mapping. CURRENT_USER and USER match the name of the current user.
PUBLIC is used to match all present and future user names in the system.

server_name

Server name of the user mapping.

OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )

Change options for the user mapping. The new options override any previously specified options. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Option names must be unique; options are also validated by the server's foreign-data wrapper.

Examples

Change the password for user mapping bob, server foo:

ALTER USER MAPPING FOR bob SERVER foo OPTIONS (SET password 'public');

Compatibility

ALTER USER MAPPING conforms to ISO/IEC 9075-9 (SQL/MED). There is a subtle syntax issue: The standard omits the FOR key word. Since both CREATE USER MAPPING and DROP USER MAPPING use FOR in analogous positions, and IBM DB2 (being the other major SQL/MED implementation) also requires it for ALTER USER MAPPING, PostgreSQL diverges from the standard here in the interest of consistency and interoperability.

See Also

CREATE USER MAPPING, DROP USER MAPPING
**ALTER VIEW**

ALTER VIEW — change the definition of a view

**Synopsis**

```
ALTER VIEW [ IF EXISTS ] name ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER VIEW [ IF EXISTS ] name ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER VIEW [ IF EXISTS ] name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER VIEW [ IF EXISTS ] name RENAME TO new_name
ALTER VIEW [ IF EXISTS ] name SET SCHEMA new_schema
ALTER VIEW [ IF EXISTS ] name SET ( view_option_name [ = view_option_value ] [, ... ] )
ALTER VIEW [ IF EXISTS ] name RESET ( view_option_name [, ... ] )
```

**Description**

ALTER VIEW changes various auxiliary properties of a view. (If you want to modify the view’s defining query, use CREATE OR REPLACE VIEW.)

You must own the view to use ALTER VIEW. To change a view’s schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the view’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the view. However, a superuser can alter ownership of any view anyway.)

**Parameters**

- **name**
  
  The name (optionally schema-qualified) of an existing view.

- **IF EXISTS**
  
  Do not throw an error if the view does not exist. A notice is issued in this case.

- **SET/DROP DEFAULT**
  
  These forms set or remove the default value for a column. A view column’s default value is substituted into any INSERT or UPDATE command whose target is the view, before applying any rules or triggers for the view. The view’s default will therefore take precedence over any default values from underlying relations.

- **new_owner**
  
  The user name of the new owner of the view.

- **new_name**
  
  The new name for the view.

- **new_schema**
  
  The new schema for the view.
ALTER VIEW

SET ( view_option_name [= view_option_value] [, ... ] )
RESET ( view_option_name [, ... ] )

Sets or resets a view option. Currently supported options are:

check_option(string)

Changes the check option of the view. The value must be local or cascaded.

security_barrier(boolean)

Changes the security-barrier property of the view. The value must be Boolean value, such as true or false.

Notes

For historical reasons, ALTER TABLE can be used with views too; but the only variants of ALTER TABLE that are allowed with views are equivalent to the ones shown above.

Examples

To rename the view foo to bar:

ALTER VIEW foo RENAME TO bar;

To attach a default column value to an updatable view:

CREATE TABLE base_table (id int, ts timestamptz);
CREATE VIEW a_view AS SELECT * FROM base_table;
ALTER VIEW a_view ALTER COLUMN ts SET DEFAULT now();
INSERT INTO base_table(id) VALUES(1);  -- ts will receive a NULL
INSERT INTO a_view(id) VALUES(2);  -- ts will receive the current time

Compatibility

ALTER VIEW is a PostgreSQL extension of the SQL standard.

See Also

CREATE VIEW, DROP VIEW
ANALYZE

ANALYZE — collect statistics about a database

Synopsis

ANALYZE [ ( option [, ...] ) ] [ table_and_columns [, ...] ]
ANALYZE [ VERBOSE ] [ table_and_columns [, ...] ]

where option can be one of:

VERBOSE

and table_and_columns is:

  table_name [ ( column_name [, ...] ) ]

Description

ANALYZE collects statistics about the contents of tables in the database, and stores the results in the pg_statistic system catalog. Subsequently, the query planner uses these statistics to help determine the most efficient execution plans for queries.

Without a table_and_columns list, ANALYZE processes every table and materialized view in the current database that the current user has permission to analyze. With a list, ANALYZE processes only those table(s). It is further possible to give a list of column names for a table, in which case only the statistics for those columns are collected.

When the option list is surrounded by parentheses, the options can be written in any order. The parenthesized syntax was added in PostgreSQL 11; the unparenthesized syntax is deprecated.

Parameters

VERBOSE

Enables display of progress messages.

table_name

The name (possibly schema-qualified) of a specific table to analyze. If omitted, all regular tables, partitioned tables, and materialized views in the current database are analyzed (but not foreign tables). If the specified table is a partitioned table, both the inheritance statistics of the partitioned table as a whole and statistics of the individual partitions are updated.

column_name

The name of a specific column to analyze. Defaults to all columns.

Outputs

When VERBOSE is specified, ANALYZE emits progress messages to indicate which table is currently being processed. Various statistics about the tables are printed as well.
Notes

Foreign tables are analyzed only when explicitly selected. Not all foreign data wrappers support **ANALYZE**. If the table's wrapper does not support **ANALYZE**, the command prints a warning and does nothing.

In the default PostgreSQL configuration, the autovacuum daemon (see Section 24.1.6) takes care of automatic analyzing of tables when they are first loaded with data, and as they change throughout regular operation. When autovacuum is disabled, it is a good idea to run **ANALYZE** periodically, or just after making major changes in the contents of a table. Accurate statistics will help the planner to choose the most appropriate query plan, and thereby improve the speed of query processing. A common strategy for read-mostly databases is to run VACUUM and **ANALYZE** once a day during a low-usage time of day. (This will not be sufficient if there is heavy update activity.)

**ANALYZE** requires only a read lock on the target table, so it can run in parallel with other activity on the table.

The statistics collected by **ANALYZE** usually include a list of some of the most common values in each column and a histogram showing the approximate data distribution in each column. One or both of these can be omitted if **ANALYZE** deems them uninteresting (for example, in a unique-key column, there are no common values) or if the column data type does not support the appropriate operators. There is more information about the statistics in Chapter 24.

For large tables, **ANALYZE** takes a random sample of the table contents, rather than examining every row. This allows even very large tables to be analyzed in a small amount of time. Note, however, that the statistics are only approximate, and will change slightly each time **ANALYZE** is run, even if the actual table contents did not change. This might result in small changes in the planner's estimated costs shown by EXPLAIN. In rare situations, this non-determinism will cause the planner's choices of query plans to change after **ANALYZE** is run. To avoid this, raise the amount of statistics collected by **ANALYZE**, as described below.

The extent of analysis can be controlled by adjusting the default_statistics_target configuration variable, or on a column-by-column basis by setting the per-column statistics target with **ALTER TABLE** ... **ALTER COLUMN** ... **SET STATISTICS** (see **ALTER TABLE**). The target value sets the maximum number of entries in the most-common-value list and the maximum number of bins in the histogram. The default target value is 100, but this can be adjusted up or down to trade off accuracy of planner estimates against the time taken for **ANALYZE** and the amount of space occupied in **pg_statistic**. In particular, setting the statistics target to zero disables collection of statistics for that column. It might be useful to do that for columns that are never used as part of the **WHERE**, **GROUP BY**, or **ORDER BY** clauses of queries, since the planner will have no use for statistics on such columns.

The largest statistics target among the columns being analyzed determines the number of table rows sampled to prepare the statistics. Increasing the target causes a proportional increase in the time and space needed to do **ANALYZE**.

One of the values estimated by **ANALYZE** is the number of distinct values that appear in each column. Because only a subset of the rows are examined, this estimate can sometimes be quite inaccurate, even with the largest possible statistics target. If this inaccuracy leads to bad query plans, a more accurate value can be determined manually and then installed with **ALTER TABLE** ... **ALTER COLUMN** ... **SET** (n_distinct = ...) (see **ALTER TABLE**).

If the table being analyzed has one or more children, **ANALYZE** will gather statistics twice: once on the rows of the parent table only, and a second time on the rows of the parent table with all of its children. This second set of statistics is needed when planning queries that traverse the entire inheritance tree. The autovacuum daemon, however, will only consider inserts or updates on the parent table itself when deciding whether to trigger an automatic analyze for that table. If that table is rarely inserted into or updated, the inheritance statistics will not be up to date unless you run **ANALYZE** manually.

If any of the child tables are foreign tables whose foreign data wrappers do not support **ANALYZE**, those child tables are ignored while gathering inheritance statistics.
If the table being analyzed is completely empty, \texttt{ANALYZE} will not record new statistics for that table. Any existing statistics will be retained.

**Compatibility**

There is no \texttt{ANALYZE} statement in the SQL standard.

**See Also**

\texttt{VACUUM}, \texttt{vacuumdb}, Section 19.4.4, Section 24.1.6
BEGIN — start a transaction block

Synopsis

BEGIN [ WORK | TRANSACTION ] [ transaction_mode [, ... ] ]

where transaction_mode is one of:

- ISOLATION LEVEL ( SERIALIZABLE | REPEATABLE READ | READ COMMITTED | READ UNCOMMITTED )
- READ WRITE | READ ONLY
- [ NOT ] DEFERRABLE

Description

BEGIN initiates a transaction block, that is, all statements after a BEGIN command will be executed in a single transaction until an explicit COMMIT or ROLLBACK is given. By default (without BEGIN), PostgreSQL executes transactions in “autocommit” mode, that is, each statement is executed in its own transaction and a commit is implicitly performed at the end of the statement (if execution was successful, otherwise a rollback is done).

Statements are executed more quickly in a transaction block, because transaction start/commit requires significant CPU and disk activity. Execution of multiple statements inside a transaction is also useful to ensure consistency when making several related changes: other sessions will be unable to see the intermediate states wherein not all the related updates have been done.

If the isolation level, read/write mode, or deferrable mode is specified, the new transaction has those characteristics, as if SET TRANSACTION was executed.

Parameters

WORK
TRANSACTION

Optional key words. They have no effect.

Refer to SET TRANSACTION for information on the meaning of the other parameters to this statement.

Notes

START TRANSACTION has the same functionality as BEGIN.

Use COMMIT or ROLLBACK to terminate a transaction block.

Issuing BEGIN when already inside a transaction block will provoke a warning message. The state of the transaction is not affected. To nest transactions within a transaction block, use savepoints (see SAVEPOINT).

For reasons of backwards compatibility, the commas between successive transaction_modes can be omitted.
Examples

To begin a transaction block:

BEGIN;

Compatibility

BEGIN is a PostgreSQL language extension. It is equivalent to the SQL-standard command START TRANSACTION, whose reference page contains additional compatibility information.

The DEFERRABLE transaction_mode is a PostgreSQL language extension.

Incidentally, the BEGIN key word is used for a different purpose in embedded SQL. You are advised to be careful about the transaction semantics when porting database applications.

See Also

COMMIT, ROLLBACK, START TRANSACTION, SAVEPOINT
CALL

CALL — invoke a procedure

Synopsis

CALL name ( [ argument ] [, ...] )

Description

CALL executes a procedure.

If the procedure has any output parameters, then a result row will be returned, containing the values of those parameters.

Parameters

name

The name (optionally schema-qualified) of the procedure.

argument

An input argument for the procedure call. See Section 4.3 for the full details on function and procedure call syntax, including use of named parameters.

Notes

The user must have EXECUTE privilege on the procedure in order to be allowed to invoke it.

To call a function (not a procedure), use SELECT instead.

If CALL is executed in a transaction block, then the called procedure cannot execute transaction control statements. Transaction control statements are only allowed if CALL is executed in its own transaction.

PL/pgSQL handles output parameters in CALL commands differently; see Section 43.6.3.

Examples

CALL do_db_maintenance();

Compatibility

CALL conforms to the SQL standard.

See Also

CREATE PROCEDURE
**CHECKPOINT**

CHECKPOINT — force a write-ahead log checkpoint

**Synopsis**

```
CHECKPOINT
```

**Description**

A checkpoint is a point in the write-ahead log sequence at which all data files have been updated to reflect the information in the log. All data files will be flushed to disk. Refer to Section 30.4 for more details about what happens during a checkpoint.

The CHECKPOINT command forces an immediate checkpoint when the command is issued, without waiting for a regular checkpoint scheduled by the system (controlled by the settings in Section 19.5.2). CHECKPOINT is not intended for use during normal operation.

If executed during recovery, the CHECKPOINT command will force a restartpoint (see Section 30.4) rather than writing a new checkpoint.

Only superusers can call CHECKPOINT.

**Compatibility**

The CHECKPOINT command is a PostgreSQL language extension.
CLOSE

CLOSE — close a cursor

Synopsis

CLOSE { name | ALL }

Description

CLOSE frees the resources associated with an open cursor. After the cursor is closed, no subsequent operations are allowed on it. A cursor should be closed when it is no longer needed.

Every non-holdable open cursor is implicitly closed when a transaction is terminated by COMMIT or ROLLBACK. A holdable cursor is implicitly closed if the transaction that created it aborts via ROLLBACK. If the creating transaction successfully commits, the holdable cursor remains open until an explicit CLOSE is executed, or the client disconnects.

Parameters

name

The name of an open cursor to close.

ALL

Close all open cursors.

Notes

PostgreSQL does not have an explicit OPEN cursor statement; a cursor is considered open when it is declared. Use the DECLARE statement to declare a cursor.

You can see all available cursors by querying the pg_cursors system view.

If a cursor is closed after a savepoint which is later rolled back, the CLOSE is not rolled back; that is, the cursor remains closed.

Examples

Close the cursor liahona:

CLOSE liahona;

Compatibility

CLOSE is fully conforming with the SQL standard. CLOSE ALL is a PostgreSQL extension.

See Also

DECLARE, FETCH, MOVE
CLUSTER

CLUSTER — cluster a table according to an index

Synopsis

CLUSTER [VERBOSE] table_name [ USING index_name ]
CLUSTER [VERBOSE]

Description

CLUSTER instructs PostgreSQL to cluster the table specified by table_name based on the index specified by index_name. The index must already have been defined on table_name.

When a table is clustered, it is physically reordered based on the index information. Clustering is a one-time operation: when the table is subsequently updated, the changes are not clustered. That is, no attempt is made to store new or updated rows according to their index order. (If one wishes, one can periodically recluster by issuing the command again. Also, setting the table's fillfactor storage parameter to less than 100% can aid in preserving cluster ordering during updates, since updated rows are kept on the same page if enough space is available there.)

When a table is clustered, PostgreSQL remembers which index it was clustered by. The form CLUSTER table_name reclusters the table using the same index as before. You can also use the CLUSTER or SET WITHOUT CLUSTER forms of ALTER TABLE to set the index to be used for future cluster operations, or to clear any previous setting.

CLUSTER without any parameter reclusters all the previously-clustered tables in the current database that the calling user owns, or all such tables if called by a superuser. This form of CLUSTER cannot be executed inside a transaction block.

When a table is being clustered, an ACCESS EXCLUSIVE lock is acquired on it. This prevents any other database operations (both reads and writes) from operating on the table until the CLUSTER is finished.

Parameters

table_name

The name (possibly schema-qualified) of a table.

index_name

The name of an index.

VERBOSE

Prints a progress report as each table is clustered.

Notes

In cases where you are accessing single rows randomly within a table, the actual order of the data in the table is unimportant. However, if you tend to access some data more than others, and there is an index that groups them together, you will benefit from using CLUSTER. If you are requesting a range of indexed values from a table, or a single indexed value that has multiple rows that match, CLUSTER will help because once the index identifies the table page for the first row that matches,
all other rows that match are probably already on the same table page, and so you save disk accesses and speed up the query.

CLUSTER can re-sort the table using either an index scan on the specified index, or (if the index is a b-tree) a sequential scan followed by sorting. It will attempt to choose the method that will be faster, based on planner cost parameters and available statistical information.

When an index scan is used, a temporary copy of the table is created that contains the table data in the index order. Temporary copies of each index on the table are created as well. Therefore, you need free space on disk at least equal to the sum of the table size and the index sizes.

When a sequential scan and sort is used, a temporary sort file is also created, so that the peak temporary space requirement is as much as double the table size, plus the index sizes. This method is often faster than the index scan method, but if the disk space requirement is intolerable, you can disable this choice by temporarily setting enable_sort to off.

It is advisable to set maintenance_work_mem to a reasonably large value (but not more than the amount of RAM you can dedicate to the CLUSTER operation) before clustering.

Because the planner records statistics about the ordering of tables, it is advisable to run ANALYZE on the newly clustered table. Otherwise, the planner might make poor choices of query plans.

Because CLUSTER remembers which indexes are clustered, one can cluster the tables one wants clustered manually the first time, then set up a periodic maintenance script that executes CLUSTER without any parameters, so that the desired tables are periodically reclustered.

Examples

Cluster the table employees on the basis of its index employees_ind:

CLUSTER employees USING employees_ind;

Cluster the employees table using the same index that was used before:

CLUSTER employees;

Cluster all tables in the database that have previously been clustered:

CLUSTER;

Compatibility

There is no CLUSTER statement in the SQL standard.

The syntax

CLUSTER index_name ON table_name

is also supported for compatibility with pre-8.3 PostgreSQL versions.

See Also

clusterdb
COMMENT

COMMENT — define or change the comment of an object

Synopsis

COMMENT ON
{
  ACCESS METHOD object_name |
  AGGREGATE aggregate_name ( aggregate_signature ) |
  CAST (source_type AS target_type) |
  COLLATION object_name |
  COLUMN relation_name.column_name |
  CONSTRAINT constraint_name ON table_name |
  CONSTRAINT constraint_name ON DOMAIN domain_name |
  CONVERSION object_name |
  DATABASE object_name |
  DOMAIN object_name |
  EXTENSION object_name |
  EVENT TRIGGER object_name |
  FOREIGN DATA WRAPPER object_name |
  FOREIGN TABLE object_name |
  FUNCTION function_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  INDEX object_name |
  LARGE OBJECT large_object_oid |
  MATERIALIZED VIEW object_name |
  OPERATOR operator_name (left_type, right_type) |
  OPERATOR CLASS object_name USING index_method |
  OPERATOR FAMILY object_name USING index_method |
  POLICY policy_name ON table_name |
  [ PROCEDURAL ] LANGUAGE object_name |
  PROCEDURE procedure_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  PUBLICATION object_name |
  ROLE object_name |
  ROUTINE routine_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  RULE rule_name ON table_name |
  SCHEMA object_name |
  SEQUENCE object_name |
  SERVER object_name |
  STATISTICS object_name |
  SUBSCRIPTION object_name |
  TABLE object_name |
  TABLESPACE object_name |
  TEXT SEARCH CONFIGURATION object_name |
  TEXT SEARCH DICTIONARY object_name |
  TEXT SEARCH PARSER object_name |
  TEXT SEARCH TEMPLATE object_name |
  TRANSFORM FOR type_name LANGUAGE lang_name |
  TRIGGER trigger_name ON table_name |
  TYPE object_name |
  VIEW object_name |
}
COMMENT

) IS 'text'

where aggregate_signature is:

+ | [ argmode ] [ argname ] argtype [ , ... ] | |
[ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ]
[ argname ] argtype [ , ... ]

Description

COMMENT stores a comment about a database object.

Only one comment string is stored for each object, so to modify a comment, issue a new COMMENT command for the same object. To remove a comment, write NULL in place of the text string. Comments are automatically dropped when their object is dropped.

For most kinds of object, only the object's owner can set the comment. Roles don't have owners, so the rule for COMMENT ON ROLE is that you must be superuser to comment on a superuser role, or have the CREATEROLE privilege to comment on non-superuser roles. Likewise, access methods don't have owners either; you must be superuser to comment on an access method. Of course, a superuser can comment on anything.

Comments can be viewed using psql's \d family of commands. Other user interfaces to retrieve comments can be built atop the same built-in functions that psql uses, namely obj_description, col_description, and shobj_description (see Table 9.68).

Parameters

object_name
relation_name.column_name
aggregate_name
constraint_name
function_name
operator_name
policy_name
procedure_name
routine_name
rule_name
trigger_name

table_name
domain_name

The name of the object to be commented. Names of tables, aggregates, collations, conversions, domains, foreign tables, functions, indexes, operators, operator classes, operator families, procedures, routines, sequences, statistics, text search objects, types, and views can be schema-qualified. When commenting on a column, relation_name must refer to a table, view, composite type, or foreign table.

table_name
domain_name

When creating a comment on a constraint, a trigger, a rule or a policy these parameters specify the name of the table or domain on which that object is defined.

source_type

The name of the source data type of the cast.

target_type

The name of the target data type of the cast.
**argmode**

The mode of a function, procedure, or aggregate argument: \texttt{IN}, \texttt{OUT}, \texttt{INOUT}, or \texttt{VARIADIC}. If omitted, the default is \texttt{IN}. Note that \texttt{COMMENT} does not actually pay any attention to \texttt{OUT} arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the \texttt{IN}, \texttt{INOUT}, and \texttt{VARIADIC} arguments.

**argname**

The name of a function, procedure, or aggregate argument. Note that \texttt{COMMENT} does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

**argtype**

The data type of a function, procedure, or aggregate argument.

**large_object_oid**

The OID of the large object.

**left_type**

**right_type**

The data type(s) of the operator’s arguments (optionally schema-qualified). Write \texttt{NONE} for the missing argument of a prefix or postfix operator.

**PROCEDURAL**

This is a noise word.

**type_name**

The name of the data type of the transform.

**lang_name**

The name of the language of the transform.

**text**

The new comment, written as a string literal; or \texttt{NULL} to drop the comment.

**Notes**

There is presently no security mechanism for viewing comments: any user connected to a database can see all the comments for objects in that database. For shared objects such as databases, roles, and tablespaces, comments are stored globally so any user connected to any database in the cluster can see all the comments for shared objects. Therefore, don’t put security-critical information in comments.

**Examples**

Attach a comment to the table \texttt{mytable}:

```
COMMENT ON TABLE mytable IS 'This is my table.';
```

Remove it again:
COMMENT ON TABLE mytable IS NULL;

Some more examples:

COMMENT ON ACCESS METHOD rtree IS 'R-Tree access method';
COMMENT ON AGGREGATE my_aggregate (double precision) IS 'Computes sample variance';
COMMENT ON CAST (text AS int4) IS 'Allow casts from text to int4';
COMMENT ON COLLATION "fr_CA" IS 'Canadian French';
COMMENT ON COLUMN my_table.my_column IS 'Employee ID number';
COMMENT ON CONVERSION my_conv IS 'Conversion to UTF8';
COMMENT ON CONSTRAINT bar_col_cons ON bar IS 'Constrains column col';
COMMENT ON CONSTRAINT dom_col_constr ON DOMAIN dom IS 'Constrains column col of domain';
COMMENT ON DATABASE my_database IS 'Development Database';
COMMENT ON DOMAIN my_domain IS 'Email Address Domain';
COMMENT ON EXTENSION hstore IS 'implements the hstore data type';
COMMENT ON FOREIGN DATA WRAPPER mywrapper IS 'my foreign data wrapper';
COMMENT ON FOREIGN TABLE my_foreign_table IS 'Employee Information in other database';
COMMENT ON FUNCTION my_function (timestamp) IS 'Returns Roman Numeral';
COMMENT ON INDEX my_index IS 'Enforces uniqueness on employee ID';
COMMENT ON LANGUAGE plpython IS 'Python support for stored procedures';
COMMENT ON LARGE OBJECT 346344 IS 'Planning document';
COMMENT ON MATERIALIZED VIEW my_matview IS 'Summary of order history';
COMMENT ON OPERATOR ^ (text, text) IS 'Performs intersection of two texts';
COMMENT ON OPERATOR - (NONE, integer) IS 'Unary minus';
COMMENT ON OPERATOR CLASS int4ops USING btree IS '4 byte integer operators for btrees';
COMMENT ON OPERATOR FAMILY integer_ops USING btree IS 'all integer operators for btrees';
COMMENT ON POLICY my_policy ON mytable IS 'Filter rows by users';
COMMENT ON PROCEDURE my_proc (integer, integer) IS 'Runs a report';
COMMENT ON ROLE my_role IS 'Administration group for finance tables';
COMMENT ON RULE my_rule ON my_table IS 'Logs updates of employee records';
COMMENT ON SCHEMA my_schema IS 'Departmental data';
COMMENT ON SEQUENCE my_sequence IS 'Used to generate primary keys';
COMMENT ON SERVER myserver IS 'my foreign server';
COMMENT ON STATISTICS my_statistics IS 'Improves planner row estimations';
COMMENT ON TABLE my_schema.my_table IS 'Employee Information';
COMMENT ON TABLESPACE my_tablespace IS 'Tablespace for indexes';
COMMENT ON TEXT SEARCH CONFIGURATION my_config IS 'Special word filtering';
COMMENT ON TEXT SEARCH DICTIONARY swedish IS 'Snowball stemmer for Swedish language';
COMMENT ON TEXT SEARCH PARSER my_parser IS 'Splits text into words';
COMMENT ON TEXT SEARCH TEMPLATE snowball IS 'Snowball stemmer';
COMMENT ON TRANSFORM FOR hstore LANGUAGE plpythonu IS 'Transform between hstore and Python dict';
COMMENT ON TRIGGER my_trigger ON my_table IS 'Used for RI';
COMMENT ON TYPE complex IS 'Complex number data type';
COMMENT ON VIEW my_view IS 'View of departmental costs';

Compatibility

There is no COMMENT command in the SQL standard.
COMMIT

COMMIT — commit the current transaction

Synopsis

COMMIT [ WORK | TRANSACTION ]

Description

COMMIT commits the current transaction. All changes made by the transaction become visible to others and are guaranteed to be durable if a crash occurs.

Parameters

WORK
TRANSACTION

Optional key words. They have no effect.

Notes

Use ROLLBACK to abort a transaction.

Issuing COMMIT when not inside a transaction does no harm, but it will provoke a warning message.

Examples

To commit the current transaction and make all changes permanent:

    COMMIT;

Compatibility

The SQL standard only specifies the two forms COMMIT and COMMIT WORK. Otherwise, this command is fully conforming.

See Also

BEGIN, ROLLBACK
COMMIT PREPARED

COMMIT PREPARED — commit a transaction that was earlier prepared for two-phase commit

Synopsis

COMMIT PREPARED transaction_id

Description

COMMIT PREPARED commits a transaction that is in prepared state.

Parameters

transaction_id

The transaction identifier of the transaction that is to be committed.

Notes

To commit a prepared transaction, you must be either the same user that executed the transaction originally, or a superuser. But you do not have to be in the same session that executed the transaction.

This command cannot be executed inside a transaction block. The prepared transaction is committed immediately.

All currently available prepared transactions are listed in the pg_prepared_xacts system view.

Examples

Commit the transaction identified by the transaction identifier foobar:

COMMIT PREPARED 'foobar';

Compatibility

COMMIT PREPARED is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.

See Also

PREPARE TRANSACTION, ROLLBACK PREPARED
COPY

COPY — copy data between a file and a table

Synopsis

COPY table_name [ ( column_name [, ...] ) ]
FROM { 'filename' | PROGRAM 'command' | STDIN }
[ [ WITH ] ( option [, ...] ) ]

COPY { table_name [ ( column_name [, ...] ) ] | ( query ) }
TO { 'filename' | PROGRAM 'command' | STDOUT }
[ [ WITH ] ( option [, ...] ) ]

where option can be one of:

- FORMAT format_name
- OIDS [ boolean ]
- FREEZE [ boolean ]
- DELIMITER 'delimiter_character'
- NULL 'null_string'
- HEADER [ boolean ]
- QUOTE 'quote_character'
- ESCAPE 'escape_character'
- FORCE_QUOTE { ( column_name [, ...] ) | * }
- FORCE_NOT_NULL ( column_name [, ...] )
- FORCE_NULL ( column_name [, ...] )
- ENCODING 'encoding_name'

Description

COPY moves data between PostgreSQL tables and standard file-system files. COPY TO copies the contents of a table to a file, while COPY FROM copies data from a file to a table (appending the data to whatever is in the table already). COPY TO can also copy the results of a SELECT query.

If a list of columns is specified, COPY will only copy the data in the specified columns to or from the file. If there are any columns in the table that are not in the column list, COPY FROM will insert the default values for those columns.

COPY with a file name instructs the PostgreSQL server to directly read from or write to a file. The file must be accessible by the PostgreSQL user (the user ID the server runs as) and the name must be specified from the viewpoint of the server. When PROGRAM is specified, the server executes the given command and reads from the standard output of the program, or writes to the standard input of the program. The command must be specified from the viewpoint of the server, and be executable by the PostgreSQL user. When stdin or stdout is specified, data is transmitted via the connection between the client and the server.

Parameters

\( table\_name \)

The name (optionally schema-qualified) of an existing table.
column_name

An optional list of columns to be copied. If no column list is specified, all columns of the table will be copied.

query

A SELECT, VALUES, INSERT, UPDATE or DELETE command whose results are to be copied. Note that parentheses are required around the query.

For INSERT, UPDATE and DELETE queries a RETURNING clause must be provided, and the target relation must not have a conditional rule, nor an ALSO rule, nor an INSTEAD rule that expands to multiple statements.

filename

The path name of the input or output file. An input file name can be an absolute or relative path, but an output file name must be an absolute path. Windows users might need to use an E'' string and double any backslashes used in the path name.

PROGRAM

A command to execute. In COPY FROM, the input is read from standard output of the command, and in COPY TO, the output is written to the standard input of the command.

Note that the command is invoked by the shell, so if you need to pass any arguments to shell command that come from an untrusted source, you must be careful to strip or escape any special characters that might have a special meaning for the shell. For security reasons, it is best to use a fixed command string, or at least avoid passing any user input in it.

STDIN

Specifies that input comes from the client application.

STDOUT

Specifies that output goes to the client application.

boolean

Specifies whether the selected option should be turned on or off. You can write TRUE, ON, or 1 to enable the option, and FALSE, OFF, or 0 to disable it. The boolean value can also be omitted, in which case TRUE is assumed.

FORMAT

Selects the data format to be read or written: text, csv (Comma Separated Values), or binary. The default is text.

OIDS

Specifies copying the OID for each row. (An error is raised if OIDS is specified for a table that does not have OIDs, or in the case of copying a query.)

FREEZE

Requests copying the data with rows already frozen, just as they would be after running the VACUUM FREEZE command. This is intended as a performance option for initial data loading. Rows will be frozen only if the table being loaded has been created or truncated in the current subtransaction, there are no cursors open and there are no older snapshots held by this transaction.
Note that all other sessions will immediately be able to see the data once it has been successfully loaded. This violates the normal rules of MVCC visibility and users specifying should be aware of the potential problems this might cause.

**DELMITER**

Specifies the character that separates columns within each row (line) of the file. The default is a tab character in text format, a comma in CSV format. This must be a single one-byte character. This option is not allowed when using binary format.

**NULL**

Specifies the string that represents a null value. The default is \N (backslash-N) in text format, and an unquoted empty string in CSV format. You might prefer an empty string even in text format for cases where you don’t want to distinguish nulls from empty strings. This option is not allowed when using binary format.

Note

When using COPY FROM, any data item that matches this string will be stored as a null value, so you should make sure that you use the same string as you used with COPY TO.

**HEADER**

Specifies that the file contains a header line with the names of each column in the file. On output, the first line contains the column names from the table, and on input, the first line is ignored. This option is allowed only when using CSV format.

**QUOTE**

Specifies the quoting character to be used when a data value is quoted. The default is double-quote. This must be a single one-byte character. This option is allowed only when using CSV format.

**ESCAPE**

Specifies the character that should appear before a data character that matches the QUOTE value. The default is the same as the QUOTE value (so that the quoting character is doubled if it appears in the data). This must be a single one-byte character. This option is allowed only when using CSV format.

**FORCE_QUOTE**

Forces quoting to be used for all non-NULL values in each specified column. NULL output is never quoted. If * is specified, non-NULL values will be quoted in all columns. This option is allowed only in COPY TO, and only when using CSV format.

**FORCE_NOT_NULL**

Do not match the specified columns' values against the null string. In the default case where the null string is empty, this means that empty values will be read as zero-length strings rather than nulls, even when they are not quoted. This option is allowed only in COPY TO, and only when using CSV format.

**FORCE_NULL**

Match the specified columns' values against the null string, even if it has been quoted, and if a match is found set the value to NULL. In the default case where the null string is empty, this converts a quoted empty string into NULL. This option is allowed only in COPY FROM, and only when using CSV format.
ENCODING

Specifies that the file is encoded in the encoding_name. If this option is omitted, the current client encoding is used. See the Notes below for more details.

Outputs

On successful completion, a COPY command returns a command tag of the form

COPY count

The count is the number of rows copied.

**Note**

psql will print this command tag only if the command was not COPY ... TO STD-OUT, or the equivalent psql meta-command \copy ... to stdout. This is to prevent confusing the command tag with the data that was just printed.

Notes

COPY TO can only be used with plain tables, not with views. However, you can write COPY (SELECT * FROM viewname) TO ... to copy the current contents of a view.

COPY FROM can be used with plain, foreign, or partitioned tables or with views that have INSTEAD OF INSERT triggers.

COPY only deals with the specific table named; it does not copy data to or from child tables. Thus for example COPY table TO shows the same data as SELECT * FROM ONLY table. But COPY (SELECT * FROM table) TO ... can be used to dump all of the data in an inheritance hierarchy.

You must have select privilege on the table whose values are read by COPY TO, and insert privilege on the table into which values are inserted by COPY FROM. It is sufficient to have column privileges on the column(s) listed in the command.

If row-level security is enabled for the table, the relevant SELECT policies will apply to COPY table TO statements. Currently, COPY FROM is not supported for tables with row-level security. Use equivalent INSERT statements instead.

Files named in a COPY command are read or written directly by the server, not by the client application. Therefore, they must reside on or be accessible to the database server machine, not the client. They must be accessible to and readable or writable by the PostgreSQL user (the user ID the server runs as), not the client. Similarly, the command specified with PROGRAM is executed directly by the server, not by the client application, must be executable by the PostgreSQL user. COPY naming a file or command is only allowed to database superusers or users who are granted one of the default roles pg_read_server_files, pg_write_server_files, or pg_execute_server_program, since it allows reading or writing any file or running a program that the server has privileges to access.

Do not confuse COPY with the psql instruction \copy. \copy invokes COPY FROM STDIN or COPY TO STDOUT, and then fetches/stores the data in a file accessible to the psql client. Thus, file accessibility and access rights depend on the client rather than the server when \copy is used.

It is recommended that the file name used in COPY always be specified as an absolute path. This is enforced by the server in the case of COPY TO, but for COPY FROM you do have the option of reading from a file specified by a relative path. The path will be interpreted relative to the working directory of the server process (normally the cluster's data directory), not the client's working directory.
Executing a command with PROGRAM might be restricted by the operating system's access control mechanisms, such as SELinux.

COPY FROM will invoke any triggers and check constraints on the destination table. However, it will not invoke rules.

For identity columns, the COPY FROM command will always write the column values provided in the input data, like the INSERT option OVERWRITING SYSTEM VALUE.

COPY input and output is affected by DateStyle. To ensure portability to other PostgreSQL installations that might use non-default DateStyle settings, DateStyle should be set to ISO before using COPY TO. It is also a good idea to avoid dumping data with IntervalStyle set to sql_standard, because negative interval values might be misinterpreted by a server that has a different setting for IntervalStyle.

Input data is interpreted according to ENCODING option or the current client encoding, and output data is encoded in ENCODING or the current client encoding, even if the data does not pass through the client but is read from or written to a file directly by the server.

COPY stops operation at the first error. This should not lead to problems in the event of a COPY TO, but the target table will already have received earlier rows in a COPY FROM. These rows will not be visible or accessible, but they still occupy disk space. This might amount to a considerable amount of wasted disk space if the failure happened well into a large copy operation. You might wish to invoke VACUUM to recover the wasted space.

FORCE_NULL and FORCE_NOT_NULL can be used simultaneously on the same column. This results in converting quoted null strings to null values and unquoted null strings to empty strings.

File Formats

Text Format

When the text format is used, the data read or written is a text file with one line per table row. Columns in a row are separated by the delimiter character. The column values themselves are strings generated by the output function, or acceptable to the input function, of each attribute's data type. The specified null string is used in place of columns that are null. COPY FROM will raise an error if any line of the input file contains more or fewer columns than are expected. If OIDS is specified, the OID is read or written as the first column, preceding the user data columns.

End of data can be represented by a single line containing just backslash-period (\.). An end-of-data marker is not necessary when reading from a file, since the end of file serves perfectly well; it is needed only when copying data to or from client applications using pre-3.0 client protocol.

Backslash characters (\) can be used in the COPY data to quote data characters that might otherwise be taken as row or column delimiters. In particular, the following characters must be preceded by a backslash if they appear as part of a column value: backslash itself, newline, carriage return, and the current delimiter character.

The specified null string is sent by COPY TO without adding any backslashes; conversely, COPY FROM matches the input against the null string before removing backslashes. Therefore, a null string such as \N cannot be confused with the actual data value \N (which would be represented as \\N).

The following special backslash sequences are recognized by COPY FROM:

<table>
<thead>
<tr>
<th>Sequence</th>
<th>Represents</th>
</tr>
</thead>
<tbody>
<tr>
<td>\b</td>
<td>Backspace (ASCII 8)</td>
</tr>
<tr>
<td>\f</td>
<td>Form feed (ASCII 12)</td>
</tr>
<tr>
<td>Sequence</td>
<td>Represents</td>
</tr>
<tr>
<td>----------</td>
<td>-----------------------------------------------</td>
</tr>
<tr>
<td>\n</td>
<td>Newline (ASCII 10)</td>
</tr>
<tr>
<td>\r</td>
<td>Carriage return (ASCII 13)</td>
</tr>
<tr>
<td>\t</td>
<td>Tab (ASCII 9)</td>
</tr>
<tr>
<td>\v</td>
<td>Vertical tab (ASCII 11)</td>
</tr>
<tr>
<td>\digits</td>
<td>Backslash followed by one to three octal digits specifies the character with that numeric code</td>
</tr>
<tr>
<td>\xdigits</td>
<td>Backslash \x followed by one or two hex digits specifies the character with that numeric code</td>
</tr>
</tbody>
</table>

Presently, "COPY TO will never emit an octal or hex-digits backslash sequence, but it does use the other sequences listed above for those control characters.

Any other backslashed character that is not mentioned in the above table will be taken to represent itself. However, beware of adding backslashes unnecessarily, since that might accidentally produce a string matching the end-of-data marker (\.) or the null string (\N by default). These strings will be recognized before any other backslash processing is done.

It is strongly recommended that applications generating "COPY data convert data newlines and carriage returns to the \n and \r sequences respectively. At present it is possible to represent a data carriage return by a backslash and carriage return, and to represent a data newline by a backslash and newline. However, these representations might not be accepted in future releases. They are also highly vulnerable to corruption if the "COPY file is transferred across different machines (for example, from Unix to Windows or vice versa).

"COPY TO will terminate each row with a Unix-style newline ("\n"). Servers running on Microsoft Windows instead output carriage return/newline ("\r\n"), but only for "COPY to a server file; for consistency across platforms, "COPY TO STDOUT always sends "\n" regardless of server platform. "COPY FROM can handle lines ending with newlines, carriage returns, or carriage return/newlines. To reduce the risk of error due to un-backslashed newlines or carriage returns that were meant as data, "COPY FROM will complain if the line endings in the input are not all alike.

**CSV Format**

This format option is used for importing and exporting the Comma Separated Value (CSV) file format used by many other programs, such as spreadsheets. Instead of the escaping rules used by PostgreSQL’s standard text format, it produces and recognizes the common CSV escaping mechanism.

The values in each record are separated by the DELIMITER character. If the value contains the delimiter character, the QUOTE character, the NULL string, a carriage return, or line feed character, then the whole value is prefixed and suffixed by the QUOTE character, and any occurrence within the value of a QUOTE character or the ESCAPE character is preceded by the escape character. You can also use FORCE_QUOTE to force quotes when outputting non-NULL values in specific columns.

The CSV format has no standard way to distinguish a NULL value from an empty string. PostgreSQL’s "COPY handles this by quoting. A NULL is output as the NULL parameter string and is not quoted, while a non-NULL value matching the NULL parameter string is quoted. For example, with the default settings, a NULL is written as an unquoted empty string, while an empty string data value is written with double quotes (" "). Reading values follows similar rules. You can use FORCE_NOT_NULL to prevent NULL input comparisons for specific columns. You can also use FORCE_NULL to convert quoted null string data values to NULL.

Because backslash is not a special character in the CSV format, \., the end-of-data marker, could also appear as a data value. To avoid any misinterpretation, a \. data value appearing as a lone entry on a line is automatically quoted on output, and on input, if quoted, is not interpreted as the end-of-data marker. If you are loading a file created by another application that has a single unquoted column and might have a value of \., you might need to quote that value in the input file.
Note

In CSV format, all characters are significant. A quoted value surrounded by white space, or any characters other than DELIMITER, will include those characters. This can cause errors if you import data from a system that pads CSV lines with white space out to some fixed width. If such a situation arises you might need to preprocess the CSV file to remove the trailing white space, before importing the data into PostgreSQL.

Note

CSV format will both recognize and produce CSV files with quoted values containing embedded carriage returns and line feeds. Thus the files are not strictly one line per table row like text-format files.

Note

Many programs produce strange and occasionally perverse CSV files, so the file format is more a convention than a standard. Thus you might encounter some files that cannot be imported using this mechanism, and COPY might produce files that other programs cannot process.

Binary Format

The binary format option causes all data to be stored/read as binary format rather than as text. It is somewhat faster than the text and CSV formats, but a binary-format file is less portable across machine architectures and PostgreSQL versions. Also, the binary format is very data type specific; for example it will not work to output binary data from a smallint column and read it into an integer column, even though that would work fine in text format.

The binary file format consists of a file header, zero or more tuples containing the row data, and a file trailer. Headers and data are in network byte order.

Note

PostgreSQL releases before 7.4 used a different binary file format.

File Header

The file header consists of 15 bytes of fixed fields, followed by a variable-length header extension area. The fixed fields are:

Signature

11-byte sequence PGCOPY\n\377\r\n\0 — note that the zero byte is a required part of the signature. (The signature is designed to allow easy identification of files that have been munged by a non-8-bit-clean transfer. This signature will be changed by end-of-line-translation filters, dropped zero bytes, dropped high bits, or parity changes.)

Flags field

32-bit integer bit mask to denote important aspects of the file format. Bits are numbered from 0 (LSB) to 31 (MSB). Note that this field is stored in network byte order (most significant byte
first), as are all the integer fields used in the file format. Bits 16-31 are reserved to denote critical
file format issues; a reader should abort if it finds an unexpected bit set in this range. Bits 0-15
are reserved to signal backwards-compatible format issues; a reader should simply ignore any
unexpected bits set in this range. Currently only one flag bit is defined, and the rest must be zero:

Bit 16

if 1, OIDs are included in the data; if 0, not

Header extension area length

32-bit integer, length in bytes of remainder of header, not including self. Currently, this is zero,
and the first tuple follows immediately. Future changes to the format might allow additional data
to be present in the header. A reader should silently skip over any header extension data it does
not know what to do with.

The header extension area is envisioned to contain a sequence of self-identifying chunks. The flags
field is not intended to tell readers what is in the extension area. Specific design of header extension
contents is left for a later release.

This design allows for both backwards-compatible header additions (add header extension chunks, or
set low-order flag bits) and non-backwards-compatible changes (set high-order flag bits to signal such
changes, and add supporting data to the extension area if needed).

Tuples

Each tuple begins with a 16-bit integer count of the number of fields in the tuple. (Presently, all tuples
in a table will have the same count, but that might not always be true.) Then, repeated for each field
in the tuple, there is a 32-bit length word followed by that many bytes of field data. (The length word
does not include itself, and can be zero.) As a special case, -1 indicates a NULL field value. No value
bytes follow in the NULL case.

There is no alignment padding or any other extra data between fields.

Presently, all data values in a binary-format file are assumed to be in binary format (format code one).
It is anticipated that a future extension might add a header field that allows per-column format codes
to be specified.

To determine the appropriate binary format for the actual tuple data you should consult the PostgreSQL
source, in particular the *send and *recv functions for each column's data type (typically these
functions are found in the src/backend/utils/adt/ directory of the source distribution).

If OIDs are included in the file, the OID field immediately follows the field-count word. It is a normal
field except that it's not included in the field-count. In particular it has a length word — this will allow
handling of 4-byte vs. 8-byte OIDs without too much pain, and will allow OIDs to be shown as null
if that ever proves desirable.

File Trailer

The file trailer consists of a 16-bit integer word containing -1. This is easily distinguished from a
tuple's field-count word.

A reader should report an error if a field-count word is neither -1 nor the expected number of columns.
This provides an extra check against somehow getting out of sync with the data.

Examples

The following example copies a table to the client using the vertical bar (|) as the field delimiter:
COPY

COPY country TO STDOUT (DELIMITER '|');

To copy data from a file into the country table:

COPY country FROM '/usr1/proj/bray/sql/country_data';

To copy into a file just the countries whose names start with 'A':

COPY (SELECT * FROM country WHERE country_name LIKE 'A%') TO '/usr1/proj/bray/sql/a_list_countries.copy';

To copy into a compressed file, you can pipe the output through an external compression program:

COPY country TO PROGRAM 'gzip > /usr1/proj/bray/sql/country_data.gz';

Here is a sample of data suitable for copying into a table from STDIN:

AF      AFGHANISTAN
AL      ALBANIA
DZ      ALGERIA
ZM      ZAMBIA
ZW      ZIMBABWE

Note that the white space on each line is actually a tab character.

The following is the same data, output in binary format. The data is shown after filtering through the Unix utility od -c. The table has three columns; the first has type char(2), the second has type text, and the third has type integer. All the rows have a null value in the third column.

| 0000000 | P | G | C | O | P | Y | n | 377 | r | n | 0 | 0 | 0 | 0 | 0 |
| 0000020 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | A |
| 0000040 | F | G | H | A | N | I | S | T | A | N | 377 | 377 | 377 | 377 | \0 |
| 0000060 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 13 |
| 0000100 | A | 377 | 377 | 377 | 377 | 377 | 377 | 377 | 377 | 377 | \0 |
| 0000120 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0000140 | Z | M | \0 | \0 | Z | A | M | B | I | A | 377 |
| 0000160 | 377 | 377 | 377 | 377 | 377 | 377 | 377 | 377 |

Compatibility

There is no COPY statement in the SQL standard.

The following syntax was used before PostgreSQL version 9.0 and is still supported:
COPY table_name [ ( column_name [, ...] ) ]
FROM { 'filename' | STDIN }
[ [ WITH ]
  [ BINARY ]
  [ OIDS ]
  [ DELIMITER [ AS ] 'delimiter' ]
  [ NULL [ AS ] 'null string' ]
  [ CSV [ HEADER ]
    [ QUOTE [ AS ] 'quote' ]
    [ ESCAPE [ AS ] 'escape' ]
    [ FORCE NOT NULL column_name [, ...] ] ] ]

COPY { table_name [ ( column_name [, ...] ) ] | ( query ) } TO { 'filename' | STDOUT }
[ [ WITH ]
  [ BINARY ]
  [ OIDS ]
  [ DELIMITER [ AS ] 'delimiter' ]
  [ NULL [ AS ] 'null string' ]
  [ CSV [ HEADER ]
    [ QUOTE [ AS ] 'quote' ]
    [ ESCAPE [ AS ] 'escape' ]
    [ FORCE QUOTE { column_name [, ...] | * } ] ] ]

Note that in this syntax, BINARY and CSV are treated as independent keywords, not as arguments of a FORMAT option.

The following syntax was used before PostgreSQL version 7.3 and is still supported:

COPY [ BINARY ] table_name [ WITH OIDS ]
FROM { 'filename' | STDIN }
[ [USING] DELIMITERS 'delimiter' ]
  [ WITH NULL AS 'null string' ]

COPY [ BINARY ] table_name [ WITH OIDS ]
TO { 'filename' | STDOUT }
[ [USING] DELIMITERS 'delimiter' ]
  [ WITH NULL AS 'null string' ]
CREATE ACCESS METHOD

CREATE ACCESS METHOD — define a new access method

Synopsis

CREATE ACCESS METHOD name
   TYPE access_method_type
   HANDLER handler_function

Description

CREATE ACCESS METHOD creates a new access method.

The access method name must be unique within the database.

Only superusers can define new access methods.

Parameters

name

The name of the access method to be created.

access_method_type

This clause specifies the type of access method to define. Only INDEX is supported at present.

handler_function

handler_function is the name (possibly schema-qualified) of a previously registered function that represents the access method. The handler function must be declared to take a single argument of type internal, and its return type depends on the type of access method; for INDEX access methods, it must be index_am_handler. The C-level API that the handler function must implement varies depending on the type of access method. The index access method API is described in Chapter 61.

Examples

Create an index access method heptree with handler function heptree_handler:

CREATE ACCESS METHOD heptree TYPE INDEX HANDLER heptree_handler;

Compatibility

CREATE ACCESS METHOD is a PostgreSQL extension.

See Also

DROP ACCESS METHOD, CREATE OPERATOR CLASS, CREATE OPERATOR FAMILY
CREATE AGGREGATE

CREATE AGGREGATE — define a new aggregate function

Synopsis

CREATE AGGREGATE name ( [ argmode ] [ argname ] arg_data_type
[ , ... ] ) ( 
    SFUNC = sfunc,
    STYPE = state_data_type
    [, SSPACE = state_data_size ]
    [, FINALFUNC = ffunc ]
    [, FINALFUNC_EXTRA ]
    [, FINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE } ]
    [, COMBINEFUNC = combinefunc ]
    [, SERIALFUNC = serialfunc ]
    [, DESERIALFUNC = deserialfunc ]
    [, INITCOND = initial_condition ]
    [, MSFUNC = msfunc ]
    [, MINVFUNC = minvfunc ]
    [, MSTYPE = mstate_data_type ]
    [, MSSPACE = mstate_data_size ]
    [, MFINALFUNC = mfunc ]
    [, MFINALFUNC_EXTRA ]
    [, MFINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE } ]
    [, MINITCOND = minitial_condition ]
    [, SORTOP = sort_operator ]
    [, PARALLEL = { SAFE | RESTRICTED | UNSAFE } ]
)

CREATE AGGREGATE name ( [ [ argmode ] [ argname ] arg_data_type
[ , ... ] ] ) ORDER BY [ argmode ] [ argname
arg_data_type [ , ... ] ) ( 
    SFUNC = sfunc,
    STYPE = state_data_type
    [, SSPACE = state_data_size ]
    [, FINALFUNC = ffunc ]
    [, FINALFUNC_EXTRA ]
    [, FINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE } ]
    [, INITCOND = initial_condition ]
    [, PARALLEL = { SAFE | RESTRICTED | UNSAFE } ]
    [, HYPOTHETICAL ]
)

or the old syntax

CREATE AGGREGATE name ( 
    BASETYPE = base_type,
    SFUNC = sfunc,
    STYPE = state_data_type
    [, SSPACE = state_data_size ]
    [, FINALFUNC = ffunc ]
)
CREATE AGGREGATE

[ , FINALFUNC_EXTRA ]
[ , FINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE } ]
[ , COMBINEFUNC = combinefunc ]
[ , SERIALFUNC = serialfunc ]
[ , DESERIALFUNC = deserialfunc ]
[ , INITCOND = initial_condition ]
[ , MSFUNC = msfunc ]
[ , MINVFUNC = minvfunc ]
[ , MSTYPE = mstate_data_type ]
[ , MSSPACE = mstate_data_size ]
[ , MFINALFUNC = mffunc ]
[ , MFINALFUNC_EXTRA ]
[ , MFINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE } ]
[ , MINITCOND = minitital_condition ]
[ , SORTOP = sort_operator ]

Description

CREATE AGGREGATE defines a new aggregate function. Some basic and commonly-used aggregate functions are included with the distribution; they are documented in Section 9.20. If one defines new types or needs an aggregate function not already provided, then CREATE AGGREGATE can be used to provide the desired features.

If a schema name is given (for example, CREATE AGGREGATE myschema.myagg ...) then the aggregate function is created in the specified schema. Otherwise it is created in the current schema.

An aggregate function is identified by its name and input data type(s). Two aggregates in the same schema can have the same name if they operate on different input types. The name and input data type(s) of an aggregate must also be distinct from the name and input data type(s) of every ordinary function in the same schema. This behavior is identical to overloading of ordinary function names (see CREATE FUNCTION).

A simple aggregate function is made from one or two ordinary functions: a state transition function sfunc, and an optional final calculation function ffunc. These are used as follows:

\[ sfunc( \text{internal-state}, \text{next-data-values} ) \longrightarrow \text{next-internal-state} \]
\[ ffunc( \text{internal-state} ) \longrightarrow \text{aggregate-value} \]

PostgreSQL creates a temporary variable of data type stype to hold the current internal state of the aggregate. At each input row, the aggregate argument value(s) are calculated and the state transition function is invoked with the current state value and the new argument value(s) to calculate a new internal state value. After all the rows have been processed, the final function is invoked once to calculate the aggregate's return value. If there is no final function then the ending state value is returned as-is.

An aggregate function can provide an initial condition, that is, an initial value for the internal state value. This is specified and stored in the database as a value of type text, but it must be a valid external representation of a constant of the state value data type. If it is not supplied then the state value starts out null.

If the state transition function is declared "strict", then it cannot be called with null inputs. With such a transition function, aggregate execution behaves as follows. Rows with any null input values are ignored (the function is not called and the previous state value is retained). If the initial state value is null, then at the first row with all-nonnull input values, the first argument value replaces the state value, and the transition function is invoked at each subsequent row with all-nonnull input values. This is handy for implementing aggregates like max. Note that this behavior is only available when
STATE DATA TYPE is the same as the first ARG DATA TYPE. WHEN THESE TYPES ARE DIFFERENT, YOU MUST SUPPLY A NONNULL INITIAL CONDITION OR USE A NONSTRICT TRANSITION FUNCTION.

IF THE STATE TRANSITION FUNCTION IS NOT STRICT, THEN IT WILL BE CALLED UNCONDITIONALLY AT EACH INPUT ROW, AND MUST DEAL WITH NULL INPUTS AND NULL STATE VALUES FOR ITSELF. THIS ALLOWS THE AGGREGATE AUTHOR TO HAVE FULL CONTROL OVER THE AGGREGATE’S HANDLING OF NULL VALUES.

IF THE FINAL FUNCTION IS DECLARED “STRICT”, THEN IT WILL NOT BE CALLED WHEN THE ENDING STATE VALUE IS NULL; INSTEAD A NULL RESULT WILL BE RETURNED AUTOMATICALLY. (OF COURSE THIS IS JUST THE NORMAL BEHAVIOR OF STRICT FUNCTIONS.) IN ANY CASE THE FINAL FUNCTION HAS THE OPTION OF RETURNING A NULL VALUE. FOR EXAMPLE, THE FINAL FUNCTION FOR AVG RETURNS NULL WHEN IT SEES THERE WERE ZERO INPUT ROWS.

SOMETIMES IT IS USEFUL TO DECLARE THE FINAL FUNCTION AS TAKING NOT JUST THE STATE VALUE, BUT EXTRA PARAMETERS CORRESPONDING TO THE AGGREGATE’S INPUT VALUES. THE MAIN REASON FOR DOING THIS IS IF THE FINAL FUNCTION IS POLYMORPHIC AND THE STATE VALUE’S DATA TYPE WOULD BE INADEQUATE TO PIN DOWN THE RESULT TYPE. THESE EXTRA PARAMETERS ARE ALWAYS PASSED AS NULL (AND SO THE FINAL FUNCTION MUST NOT BE STRICT WHEN THE FINALFUNC_EXTRA OPTION IS USED), BUT NONETHLESS THEY ARE VALID PARAMETERS. THE FINAL FUNCTION COULD FOR EXAMPLE MAKE USE OF GET_FN_EXPRARGTYPE TO IDENTIFY THE ACTUAL ARGUMENT TYPE IN THE CURRENT CALL.

AN AGGREGATE CAN Optionally SUPPORT MOVING-AGGREGATE MODE, AS DESCRIBED IN SECTION 38.11.1. THIS REQUIRES SPECIFYING THE MSFUNC, MINVFUNC, AND MSTYPE PARAMETERS, AND Optionally THE MS-PACE, MFINALFUNC, MFINALFUNC_EXTRA, MFINALFUNC_MODIFY, AND MINITCOND PARAMETERS. EXCEPT FOR MINVFUNC, THESE PARAMETERS WORK LIKE THE CORRESPONDING SIMPLE-AGGREGATE PARAMETERS WITHOUT M; THEY DEFINE A SEPARATE IMPLEMENTATION OF THE AGGREGATE THAT INCLUDES AN INVERSE TRANSITION FUNCTION.

THE SYNTAX WITH ORDER BY IN THE PARAMETER LIST CREATES A SPECIAL TYPE OF AGGREGATE CALLED AN ORDERED-SET AGGREGATE; OR IF HYPOTHETICAL IS SPECIFIED, THEN A HYPOTHETICAL-SET AGGREGATE IS CREATED. THESE AGGREGATES OPERATE OVER GROUPS OF SORTED VALUES IN ORDER-DEPENDENT WAYS, SO THAT SPECIFICATION OF AN INPUT SORT ORDER IS AN ESSENTIAL PART OF A CALL. ALSO, THEY CAN HAVE DIRECT ARGUMENTS, WHICH ARE ARGUMENTS THAT ARE EVALUATED ONLY ONCE PER AGGREGATION RATHER THAN ONCE PER INPUT ROW. HYPOTHETICAL-SET AGGREGATES ARE A SUBCLASS OF ORDERED-SET AGGREGATES IN WHICH SOME OF THE DIRECT ARGUMENTS ARE REQUIRED TO MATCH, IN NUMBER AND DATA TYPES, THE AGGREGATED ARGUMENT COLUMNS. THIS ALLOW IN THE VALUES OF THOSE DIRECT ARGUMENTS TO BE ADDED TO THE COLLECTION OF AGGREGATE-INPUT ROWS AS AN ADDITIONAL “HYPOTHETICAL” ROW.

AN AGGREGATE CAN Optionally SUPPORT PARTIAL AGGREGATION, AS DESCRIBED IN SECTION 38.11.4. THIS REQUIRES SPECIFYING THE COMBINEFUNC PARAMETER. IF THE STATE DATA TYPE IS INTERNAL, IT’S Usually ALSO APPROPRIATE TO PROVIDE THE SERIALFUNC AND DEserializer PARAMETERS SO THAT PARALLEL AGGREGATION IS POSSIBLE. NOTE THAT THE AGGREGATE MUST ALSO BE MARKED PARALLEL SAFE TO ENABLE PARALLEL AGGREGATION.

AGGREGATES THAT BEHAVE LIKE MIN OR MAX CAN SOMETIMES BE OPTIMIZED BY LOOKING INTO AN INDEX INSTEAD OF SCANNING EVERY INPUT ROW. IF THIS AGGREGATE CAN BE SO OPTIMIZED, INDICATE IT BY SPECIFYING A SORT OPERATOR. THE BASIC REQUIREMENT IS THAT THE AGGREGATE MUST YIELD THE FIRST ELEMENT IN THE SORT ORDERING INDUCED BY THE OPERATOR; IN OTHER WORDS:

SELECT agg(col) FROM tab;

MUST BE EQUIVALENT TO:

SELECT col FROM tab ORDER BY col USING sortop LIMIT 1;

FURTHER ASSUMPTIONS ARE THAT THE AGGREGATE IGNORES NULL INPUTS, AND THAT IT DELIVERS A NULL RESULT IF AND ONLY IF THERE WERE NO NON-NULL INPUTS. ORDINARILY, A DATA TYPE’S < OPERATOR IS THE PROPER SORT OPERATOR FOR MIN, AND > IS THE PROPER SORT OPERATOR FOR MAX. NOTE THAT THE OPTIMIZATION WILL NEVER ACTUALLY TAKE EFFECT UNLESS THE SPECIFIED OPERATOR IS THE “LESS THAN” OR “GREATER THAN” STRATEGY MEMBER OF A B-TREE INDEX OPERATOR CLASS.
CREATE AGGREGATE

To be able to create an aggregate function, you must have USAGE privilege on the argument types, the state type(s), and the return type, as well as EXECUTE privilege on the supporting functions.

Parameters

name

The name (optionally schema-qualified) of the aggregate function to create.

argmode

The mode of an argument: IN or VARIADIC. (Aggregate functions do not support OUT arguments.) If omitted, the default is IN. Only the last argument can be marked VARIADIC.

argname

The name of an argument. This is currently only useful for documentation purposes. If omitted, the argument has no name.

arg_data_type

An input data type on which this aggregate function operates. To create a zero-argument aggregate function, write * in place of the list of argument specifications. (An example of such an aggregate is \texttt{count(*)}.)

base_type

In the old syntax for \texttt{CREATE AGGREGATE}, the input data type is specified by a \texttt{basetype} parameter rather than being written next to the aggregate name. Note that this syntax allows only one input parameter. To define a zero-argument aggregate function with this syntax, specify the \texttt{basetype} as "ANY" (not *). Ordered-set aggregates cannot be defined with the old syntax.

sfunc

The name of the state transition function to be called for each input row. For a normal $N$-argument aggregate function, the \texttt{sfunc} must take $N+1$ arguments, the first being of type \texttt{state_data_type} and the rest matching the declared input data type(s) of the aggregate. The function must return a value of type \texttt{state_data_type}. This function takes the current state value and the current input data value(s), and returns the next state value.

For ordered-set (including hypothetical-set) aggregates, the state transition function receives only the current state value and the aggregated arguments, not the direct arguments. Otherwise it is the same.

state_data_type

The data type for the aggregate's state value.

state_data_size

The approximate average size (in bytes) of the aggregate's state value. If this parameter is omitted or is zero, a default estimate is used based on the \texttt{state_data_type}. The planner uses this value to estimate the memory required for a grouped aggregate query. The planner will consider using hash aggregation for such a query only if the hash table is estimated to fit in work_mem; therefore, large values of this parameter discourage use of hash aggregation.

ffunc

The name of the final function called to compute the aggregate's result after all input rows have been traversed. For a normal aggregate, this function must take a single argument of type \texttt{state_data_type}. The return data type of the aggregate is defined as the return type of this function. If \texttt{ffunc} is not specified, then the ending state value is used as the aggregate's result, and the return type is \texttt{state_data_type}. 

1490
For ordered-set (including hypothetical-set) aggregates, the final function receives not only the final state value, but also the values of all the direct arguments.

If FINALFUNC_EXTRA is specified, then in addition to the final state value and any direct arguments, the final function receives extra NULL values corresponding to the aggregate's regular (aggregated) arguments. This is mainly useful to allow correct resolution of the aggregate result type when a polymorphic aggregate is being defined.

FINALFUNC_MODIFY = { READ_ONLY | SHAREABLE | READ_WRITE }

This option specifies whether the final function is a pure function that does not modify its arguments. READ_ONLY indicates it does not; the other two values indicate that it may change the transition state value. See Notes below for more detail. The default is READ_ONLY, except for ordered-set aggregates, for which the default is READ_WRITE.

combinefunc

The combinefunc function may optionally be specified to allow the aggregate function to support partial aggregation. If provided, the combinefunc must combine two state_data_type values, each containing the result of aggregation over some subset of the input values, to produce a new state_data_type that represents the result of aggregating over both sets of inputs. This function can be thought of as an sfunc, where instead of acting upon an individual input row and adding it to the running aggregate state, it adds another aggregate state to the running state.

The combinefunc must be declared as taking two arguments of the state_data_type and returning a value of the state_data_type. Optionally this function may be “strict”. In this case the function will not be called when either of the input states are null; the other state will be taken as the correct result.

For aggregate functions whose state_data_type is internal, the combinefunc must not be strict. In this case the combinefunc must ensure that null states are handled correctly and that the state being returned is properly stored in the aggregate memory context.

serialfunc

An aggregate function whose state_data_type is internal can participate in parallel aggregation only if it has a serialfunc function, which must serialize the aggregate state into a bytea value for transmission to another process. This function must take a single argument of type internal and return type bytea. A corresponding deserialfunc is also required.

deserialfunc

Deserializes a previously serialized aggregate state back into state_data_type. This function must take two arguments of types bytea and internal, and produce a result of type internal. (Note: the second, internal argument is unused, but is required for type safety reasons.)

initial_condition

The initial setting for the state value. This must be a string constant in the form accepted for the data type state_data_type. If not specified, the state value starts out null.

msfunc

The name of the forward state transition function to be called for each input row in moving-aggregate mode. This is exactly like the regular transition function, except that its first argument and result are of type mstate_data_type, which might be different from state_data_type.

minvfunc

The name of the inverse state transition function to be used in moving-aggregate mode. This function has the same argument and result types as msfunc, but it is used to remove a value from
the current aggregate state, rather than add a value to it. The inverse transition function must have
the same strictness attribute as the forward state transition function.

\textbf{mstate\_data\_type}

The data type for the aggregate's state value, when using moving-aggregate mode.

\textbf{mstate\_data\_size}

The approximate average size (in bytes) of the aggregate's state value, when using moving-ag-
gregate mode. This works the same as \textit{state\_data\_size}.

\textbf{mffunc}

The name of the final function called to compute the aggregate's result after all input rows have
been traversed, when using moving-aggregate mode. This works the same as \textit{ffunc}, except
that its first argument's type is \textit{mstate\_data\_type} and extra dummy arguments are specified by writing \texttt{MFINALFUNC\_EXTRA}. The aggregate result type determined by \textit{mffunc} or
\textit{mstate\_data\_type} must match that determined by the aggregate's regular implementation.

\texttt{MFINALFUNC\_MODIFY = \{ READ\_ONLY | SHAREABLE | READ\_WRITE \}}

This option is like \texttt{FINALFUNC\_MODIFY}, but it describes the behavior of the moving-aggregate
final function.

\textbf{minital\_condition}

The initial setting for the state value, when using moving-aggregate mode. This works the same
as \textit{initial\_condition}.

\textbf{sort\_operator}

The associated sort operator for a \texttt{MIN} or \texttt{MAX}-like aggregate. This is just an operator name
(possibly schema-qualified). The operator is assumed to have the same input data types as the
aggregate (which must be a single-argument normal aggregate).

\texttt{PARALLEL = \{ SAFE | RESTRICTED | UNSAFE \}}

The meanings of \texttt{PARALLEL SAFE}, \texttt{PARALLEL RESTRICTED}, and \texttt{PARALLEL UNSAFE} are
the same as in \texttt{CREATE FUNCTION}. An aggregate will not be considered for parallelization if it
is marked \texttt{PARALLEL UNSAFE} (which is the default!) or \texttt{PARALLEL RESTRICTED}. Note that
the parallel-safety markings of the aggregate's support functions are not consulted by the planner,
only the marking of the aggregate itself.

\textbf{HYPOTHETICAL}

For ordered-set aggregates only, this flag specifies that the aggregate arguments are to be
processed according to the requirements for hypothetical-set aggregates: that is, the last few di-
rect arguments must match the data types of the aggregated (\texttt{WITHIN GROUP}) arguments. The
\textbf{HYPOTHETICAL} flag has no effect on run-time behavior, only on parse-time resolution of the
data types and collations of the aggregate's arguments.

The parameters of \texttt{CREATE AGGREGATE} can be written in any order, not just the order illustrated
above.

\textbf{Notes}

In parameters that specify support function names, you can write a schema name if needed, for example
\texttt{SFUNC = public.sum}. Do not write argument types there, however — the argument types of the
support functions are determined from other parameters.
Ordinarily, PostgreSQL functions are expected to be true functions that do not modify their input values. However, an aggregate transition function, when used in the context of an aggregate, is allowed to cheat and modify its transition-state argument in place. This can provide substantial performance benefits compared to making a fresh copy of the transition state each time.

Likewise, while an aggregate final function is normally expected not to modify its input values, sometimes it is impractical to avoid modifying the transition-state argument. Such behavior must be declared using the FINALFUNC_MODIFY parameter. The READ_WRITE value indicates that the final function modifies the transition state in unspecified ways. This value prevents use of the aggregate as a window function, and it also prevents merging of transition states for aggregate calls that share the same input values and transition functions. The SHAREABLE value indicates that the transition function cannot be applied after the final function, but multiple final-function calls can be performed on the ending transition state value. This value prevents use of the aggregate as a window function, but it allows merging of transition states. (That is, the optimization of interest here is not applying the same final function repeatedly, but applying different final functions to the same ending transition state value. This is allowed as long as none of the final functions are marked READ_WRITE.)

If an aggregate supports moving-aggregate mode, it will improve calculation efficiency when the aggregate is used as a window function for a window with moving frame start (that is, a frame start mode other than UNBOUNDED PRECEDING). Conceptually, the forward transition function adds input values to the aggregate's state when they enter the window frame from the bottom, and the inverse transition function removes them again when they leave the frame at the top. So, when values are removed, they are always removed in the same order they were added. Whenever the inverse transition function is invoked, it will thus receive the earliest added but not yet removed argument value(s). The inverse transition function can assume that at least one row will remain in the current state after it removes the oldest row. (When this would not be the case, the window function mechanism simply starts a fresh aggregation, rather than using the inverse transition function.)

The forward transition function for moving-aggregate mode is not allowed to return NULL as the new state value. If the inverse transition function returns NULL, this is taken as an indication that the inverse function cannot reverse the state calculation for this particular input, and so the aggregate calculation will be redone from scratch for the current frame starting position. This convention allows moving-aggregate mode to be used in situations where there are some infrequent cases that are impractical to reverse out of the running state value.

If no moving-aggregate implementation is supplied, the aggregate can still be used with moving frames, but PostgreSQL will recompute the whole aggregation whenever the start of the frame moves. Note that whether or not the aggregate supports moving-aggregate mode, PostgreSQL can handle a moving frame end without recalculation; this is done by continuing to add new values to the aggregate's state. This is why use of an aggregate as a window function requires that the final function be read-only: it must not damage the aggregate's state value, so that the aggregation can be continued even after an aggregate result value has been obtained for one set of frame boundaries.

The syntax for ordered-set aggregates allows VARIADIC to be specified for both the last direct parameter and the last aggregated parameter. However, the current implementation restricts use of VARIADIC in two ways. First, ordered-set aggregates can only use VARIADIC "any", not other variadic array types. Second, if the last direct parameter is VARIADIC "any", then there can be only one aggregated parameter and it must also be VARIADIC "any". (In the representation used in the system catalogs, these two parameters are merged into a single VARIADIC "any" item, since pg_proc cannot represent functions with more than one VARIADIC parameter.) If the aggregate is a hypothetical-set aggregate, the direct arguments that match the VARIADIC "any" parameter are the hypothetical ones; any preceding parameters represent additional direct arguments that are not constrained to match the aggregated arguments.

Currently, ordered-set aggregates do not need to support moving-aggregate mode, since they cannot be used as window functions.

Partial (including parallel) aggregation is currently not supported for ordered-set aggregates. Also, it will never be used for aggregate calls that include DISTINCT or ORDER BY clauses, since those semantics cannot be supported during partial aggregation.
CREATE AGGREGATE

Examples

See Section 38.11.

Compatibility

CREATE AGGREGATE is a PostgreSQL language extension. The SQL standard does not provide for user-defined aggregate functions.

See Also

ALTER AGGREGATE, DROP AGGREGATE
CREATE CAST

CREATE CAST — define a new cast

Synopsis

CREATE CAST (source_type AS target_type)
   WITH FUNCTION function_name [ (argument_type [, ...]) ]
   [ AS ASSIGNMENT | AS IMPLICIT ]

CREATE CAST (source_type AS target_type)
   WITHOUT FUNCTION
   [ AS ASSIGNMENT | AS IMPLICIT ]

CREATE CAST (source_type AS target_type)
   WITH INOUT
   [ AS ASSIGNMENT | AS IMPLICIT ]

Description

CREATE CAST defines a new cast. A cast specifies how to perform a conversion between two data types. For example,

SELECT CAST(42 AS float8);

converts the integer constant 42 to type float8 by invoking a previously specified function, in this case float8(int4). (If no suitable cast has been defined, the conversion fails.)

Two types can be binary coercible, which means that the conversion can be performed “for free” without invoking any function. This requires that corresponding values use the same internal representation. For instance, the types text and varchar are binary coercible both ways. Binary coercibility is not necessarily a symmetric relationship. For example, the cast from xml to text can be performed for free in the present implementation, but the reverse direction requires a function that performs at least a syntax check. (Two types that are binary coercible both ways are also referred to as binary compatible.)

You can define a cast as an I/O conversion cast by using the WITH INOUT syntax. An I/O conversion cast is performed by invoking the output function of the source data type, and passing the resulting string to the input function of the target data type. In many common cases, this feature avoids the need to write a separate cast function for conversion. An I/O conversion cast acts the same as a regular function-based cast; only the implementation is different.

By default, a cast can be invoked only by an explicit cast request, that is an explicit CAST(x AS typename) or x::typename construct.

If the cast is marked AS ASSIGNMENT then it can be invoked implicitly when assigning a value to a column of the target data type. For example, supposing that foo.f1 is a column of type text, then:

INSERT INTO foo (f1) VALUES (42);

will be allowed if the cast from type integer to type text is marked AS ASSIGNMENT, otherwise not. (We generally use the term assignment cast to describe this kind of cast.)
CREATE CAST

If the cast is marked AS IMPLICIT then it can be invoked implicitly in any context, whether assignment or internally in an expression. (We generally use the term implicit cast to describe this kind of cast.) For example, consider this query:

SELECT 2 + 4.0;

The parser initially marks the constants as being of type integer and numeric respectively. There is no integer + numeric operator in the system catalogs, but there is a numeric + numeric operator. The query will therefore succeed if a cast from integer to numeric is available and is marked AS IMPLICIT — which in fact it is. The parser will apply the implicit cast and resolve the query as if it had been written

SELECT CAST ( 2 AS numeric ) + 4.0;

Now, the catalogs also provide a cast from numeric to integer. If that cast were marked AS IMPLICIT — which it is not — then the parser would be faced with choosing between the above interpretation and the alternative of casting the numeric constant to integer and applying the integer + integer operator. Lacking any knowledge of which choice to prefer, it would give up and declare the query ambiguous. The fact that only one of the two casts is implicit is the way in which we teach the parser to prefer resolution of a mixed numeric-and-integer expression as numeric; there is no built-in knowledge about that.

It is wise to be conservative about marking casts as implicit. An overabundance of implicit casting paths can cause PostgreSQL to choose surprising interpretations of commands, or to be unable to resolve commands at all because there are multiple possible interpretations. A good rule of thumb is to make a cast implicitly invokable only for information-preserving transformations between types in the same general type category. For example, the cast from int2 to int4 can reasonably be implicit, but the cast from float8 to int4 should probably be assignment-only. Cross-type-category casts, such as text to int4, are best made explicit-only.

**Note**

Sometimes it is necessary for usability or standards-compliance reasons to provide multiple implicit casts among a set of types, resulting in ambiguity that cannot be avoided as above. The parser has a fallback heuristic based on type categories and preferred types that can help to provide desired behavior in such cases. See CREATE TYPE for more information.

To be able to create a cast, you must own the source or the target data type and have USAGE privilege on the other type. To create a binary-coercible cast, you must be superuser. (This restriction is made because an erroneous binary-coercible cast conversion can easily crash the server.)

**Parameters**

- **source_type**
  
  The name of the source data type of the cast.

- **target_type**
  
  The name of the target data type of the cast.

- **function_name**(argument_type [, ...])
  
  The function used to perform the cast. The function name can be schema-qualified. If it is not, the function will be looked up in the schema search path. The function's result data type must match
CREATE CAST

the target type of the cast. Its arguments are discussed below. If no argument list is specified, the
function name must be unique in its schema.

WITHOUT FUNCTION

Indicates that the source type is binary-coercible to the target type, so no function is required to
perform the cast.

WITH INOUT

Indicates that the cast is an I/O conversion cast, performed by invoking the output function of the
source data type, and passing the resulting string to the input function of the target data type.

AS ASSIGNMENT

Indicates that the cast can be invoked implicitly in assignment contexts.

AS IMPLICIT

Indicates that the cast can be invoked implicitly in any context.

Cast implementation functions can have one to three arguments. The first argument type must be
identical to or binary-coercible from the cast's source type. The second argument, if present, must be
type integer; it receives the type modifier associated with the destination type, or -1 if there is
none. The third argument, if present, must be type boolean; it receives true if the cast is an explicit
cast, false otherwise. (Bizarrely, the SQL standard demands different behaviors for explicit and
implicit casts in some cases. This argument is supplied for functions that must implement such casts.
It is not recommended that you design your own data types so that this matters.)

The return type of a cast function must be identical to or binary-coercible to the cast's target type.

Ordinarily a cast must have different source and target data types. However, it is allowed to declare
a cast with identical source and target types if it has a cast implementation function with more than
one argument. This is used to represent type-specific length coercion functions in the system catalogs.
The named function is used to coerce a value of the type to the type modifier value given by its second
argument.

When a cast has different source and target types and a function that takes more than one argument,
it supports converting from one type to another and applying a length coercion in a single step. When
no such entry is available, coercion to a type that uses a type modifier involves two cast steps, one to
convert between data types and a second to apply the modifier.

A cast to or from a domain type currently has no effect. Casting to or from a domain uses the casts
associated with its underlying type.

Notes

Use DROP CAST to remove user-defined casts.

Remember that if you want to be able to convert types both ways you need to declare casts both ways
explicitly.

It is normally not necessary to create casts between user-defined types and the standard string types
(text, varchar, and char (n), as well as user-defined types that are defined to be in the string
category). PostgreSQL provides automatic I/O conversion casts for that. The automatic casts to string
types are treated as assignment casts, while the automatic casts from string types are explicit-only.
You can override this behavior by declaring your own cast to replace an automatic cast, but usually
the only reason to do so is if you want the conversion to be more easily invokable than the standard
assignment-only or explicit-only setting. Another possible reason is that you want the conversion to
behave differently from the type's I/O function; but that is sufficiently surprising that you should think twice about whether it's a good idea. (A small number of the built-in types do indeed have different behaviors for conversions, mostly because of requirements of the SQL standard.)

While not required, it is recommended that you continue to follow this old convention of naming cast implementation functions after the target data type. Many users are used to being able to cast data types using a function-style notation, that is `typename(x)`. This notation is in fact nothing more nor less than a call of the cast implementation function; it is not specially treated as a cast. If your conversion functions are not named to support this convention then you will have surprised users. Since PostgreSQL allows overloading of the same function name with different argument types, there is no difficulty in having multiple conversion functions from different types that all use the target type's name.

```
Note
Actually the preceding paragraph is an oversimplification: there are two cases in which a function-call construct will be treated as a cast request without having matched it to an actual function. If a function call `name(x)` does not exactly match any existing function, but `name` is the name of a data type and `pg_cast` provides a binary-coercible cast to this type from the type of `x`, then the call will be construed as a binary-coercible cast. This exception is made so that binary-coercible casts can be invoked using functional syntax, even though they lack any function. Likewise, if there is no `pg_cast` entry but the cast would be to or from a string type, the call will be construed as an I/O conversion cast. This exception allows I/O conversion casts to be invoked using functional syntax.
```

```
Note
There is also an exception to the exception: I/O conversion casts from composite types to string types cannot be invoked using functional syntax, but must be written in explicit cast syntax (either `CAST` or `::` notation). This exception was added because after the introduction of automatically-provided I/O conversion casts, it was found too easy to accidentally invoke such a cast when a function or column reference was intended.
```

Examples

To create an assignment cast from type `bigint` to type `int4` using the function `int4(bigint)`:  

```
CREATE CAST (bigint AS int4) WITH FUNCTION int4(bigint) AS ASSIGNMENT;
```

(This cast is already predefined in the system.)

Compatibility

The `CREATE CAST` command conforms to the SQL standard, except that SQL does not make provisions for binary-coercible types or extra arguments to implementation functions. `AS IMPLICIT` is a PostgreSQL extension, too.

See Also

CREATE FUNCTION, CREATE TYPE, DROP CAST
CREATE COLLATION

CREATE COLLATION — define a new collation

Synopsis

```
CREATE COLLATION [ IF NOT EXISTS ] name ( 
   [ LOCALE = locale, ]
   [ LC_COLLATE = lc_collate, ]
   [ LC_CTYPE = lc_ctype, ]
   [ PROVIDER = provider, ]
   [ VERSION = version ]
)
CREATE COLLATION [ IF NOT EXISTS ] name FROM existing_collation
```

Description

CREATE_COLLATION defines a new collation using the specified operating system locale settings, or by copying an existing collation.

To be able to create a collation, you must have CREATE privilege on the destination schema.

Parameters

**IF NOT EXISTS**

Do not throw an error if a collation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing collation is anything like the one that would have been created.

**name**

The name of the collation. The collation name can be schema-qualified. If it is not, the collation is defined in the current schema. The collation name must be unique within that schema. (The system catalogs can contain collations with the same name for other encodings, but these are ignored if the database encoding does not match.)

**locale**

This is a shortcut for setting LC_COLLATE and LC_CTYPE at once. If you specify this, you cannot specify either of those parameters.

**lc_collate**

Use the specified operating system locale for the LC_COLLATE locale category.

**lc_ctype**

Use the specified operating system locale for the LC_CTYPE locale category.

**provider**

Specifies the provider to use for locale services associated with this collation. Possible values are: icu, libc. libc is the default. The available choices depend on the operating system and build options.
version

Specifies the version string to store with the collation. Normally, this should be omitted, which will cause the version to be computed from the actual version of the collation as provided by the operating system. This option is intended to be used by `pg_upgrade` for copying the version from an existing installation.

See also ALTER COLLATION for how to handle collation version mismatches.

existing_collation

The name of an existing collation to copy. The new collation will have the same properties as the existing one, but it will be an independent object.

Notes

Use DROP COLLATION to remove user-defined collations.

See Section 23.2.2.3 for more information on how to create collations.

When using the libc collation provider, the locale must be applicable to the current database encoding. See CREATE DATABASE for the precise rules.

Examples

To create a collation from the operating system locale `fr_FR.utf8` (assuming the current database encoding is UTF8):

```
CREATE COLLATION french (locale = 'fr_FR.utf8');
```

To create a collation using the ICU provider using German phone book sort order:

```
CREATE COLLATION german_phonebook (provider = icu, locale = 'de-u-co-phonebk');
```

To create a collation from an existing collation:

```
CREATE COLLATION german FROM "de_DE";
```

This can be convenient to be able to use operating-system-independent collation names in applications.

Compatibility

There is a CREATE COLLATION statement in the SQL standard, but it is limited to copying an existing collation. The syntax to create a new collation is a PostgreSQL extension.

See Also

ALTER COLLATION, DROP COLLATION
CREATE CONVERSION

CREATE CONVERSION — define a new encoding conversion

Synopsis

CREATE [ DEFAULT ] CONVERSION name
   FOR source_encoding TO dest_encoding FROM function_name

Description

CREATE CONVERSION defines a new conversion between character set encodings. Also, conversions that are marked DEFAULT can be used for automatic encoding conversion between client and server. For this purpose, two conversions, from encoding A to B and from encoding B to A, must be defined.

To be able to create a conversion, you must have EXECUTE privilege on the function and CREATE privilege on the destination schema.

Parameters

DEFAULT

The DEFAULT clause indicates that this conversion is the default for this particular source to destination encoding. There should be only one default encoding in a schema for the encoding pair.

name

The name of the conversion. The conversion name can be schema-qualified. If it is not, the conversion is defined in the current schema. The conversion name must be unique within a schema.

source_encoding

The source encoding name.

dest_encoding

The destination encoding name.

function_name

The function used to perform the conversion. The function name can be schema-qualified. If it is not, the function will be looked up in the path.

The function must have the following signature:

conv_proc(
   integer,  -- source encoding ID
   integer,  -- destination encoding ID
   cstring,  -- source string (null terminated C string)
   internal, -- destination (fill with a null terminated C string)
   integer   -- source string length
) RETURNS void;
Notes

Use \texttt{DROP CONVERSION} to remove user-defined conversions.

The privileges required to create a conversion might be changed in a future release.

Examples

To create a conversion from encoding \texttt{UTF8} to \texttt{LATIN1} using \texttt{myfunc}:

\begin{verbatim}
CREATE CONVERSION myconv FOR 'UTF8' TO 'LATIN1' FROM myfunc;
\end{verbatim}

Compatibility

\texttt{CREATE CONVERSION} is a PostgreSQL extension. There is no \texttt{CREATE CONVERSION} statement in the SQL standard, but a \texttt{CREATE TRANSLATION} statement that is very similar in purpose and syntax.

See Also

\texttt{ALTER CONVERSION}, \texttt{CREATE FUNCTION}, \texttt{DROP CONVERSION}
CREATE DATABASE

CREATE DATABASE — create a new database

Synopsis

CREATE DATABASE name
    [ [ WITH ] [ OWNER [=] user_name ]
    [ TEMPLATE [=] template ]
    [ ENCODING [=] encoding ]
    [ LC_COLLATE [=] lc_collate ]
    [ LC_CTYPE [=] lc_ctype ]
    [ TABLESPACE [=] tablespace_name ]
    [ ALLOW_CONNECTIONS [=] allowconn ]
    [ CONNECTION LIMIT [=] connlimit ]
    [ IS_TEMPLATE [=] istemplate ]

Description

CREATE DATABASE creates a new PostgreSQL database.

To create a database, you must be a superuser or have the special CREATEDB privilege. See CREATE ROLE.

By default, the new database will be created by cloning the standard system database template1. A different template can be specified by writing TEMPLATE name. In particular, by writing TEMPLATE template0, you can create a virgin database containing only the standard objects predefined by your version of PostgreSQL. This is useful if you wish to avoid copying any installation-local objects that might have been added to template1.

Parameters

name

The name of a database to create.

user_name

The role name of the user who will own the new database, or DEFAULT to use the default (namely, the user executing the command). To create a database owned by another role, you must be a direct or indirect member of that role, or be a superuser.

template

The name of the template from which to create the new database, or DEFAULT to use the default template (template1).

encoding

Character set encoding to use in the new database. Specify a string constant (e.g., 'SQL_ASCII'), or an integer encoding number, or DEFAULT to use the default encoding (namely, the encoding of the template database). The character sets supported by the PostgreSQL server are described in Section 23.3.1. See below for additional restrictions.
CREATE DATABASE

lc_collate

Collation order (LC_COLLATE) to use in the new database. This affects the sort order applied to strings, e.g. in queries with ORDER BY, as well as the order used in indexes on text columns. The default is to use the collation order of the template database. See below for additional restrictions.

lc_ctype

Character classification (LC_CTYPE) to use in the new database. This affects the categorization of characters, e.g. lower, upper and digit. The default is to use the character classification of the template database. See below for additional restrictions.

tablespace_name

The name of the tablespace that will be associated with the new database, or DEFAULT to use the template database's tablespace. This tablespace will be the default tablespace used for objects created in this database. See CREATE TABLESPACE for more information.

allowconn

If false then no one can connect to this database. The default is true, allowing connections (except as restricted by other mechanisms, such as GRANT/REVOKE CONNECT).

conllimit

How many concurrent connections can be made to this database. -1 (the default) means no limit.

istemplate

If true, then this database can be cloned by any user with CREATEDB privileges; if false (the default), then only superusers or the owner of the database can clone it.

Optional parameters can be written in any order, not only the order illustrated above.

Notes

CREATE DATABASE cannot be executed inside a transaction block.

Errors along the line of “could not initialize database directory” are most likely related to insufficient permissions on the data directory, a full disk, or other file system problems.

Use DROP DATABASE to remove a database.

The program createdb is a wrapper program around this command, provided for convenience.

Database-level configuration parameters (set via ALTER DATABASE) are not copied from the template database.

Although it is possible to copy a database other than template1 by specifying its name as the template, this is not (yet) intended as a general-purpose “COPY DATABASE” facility. The principal limitation is that no other sessions can be connected to the template database while it is being copied. CREATE DATABASE will fail if any other connection exists when it starts; otherwise, new connections to the template database are locked out until CREATE DATABASE completes. See Section 22.3 for more information.

The character set encoding specified for the new database must be compatible with the chosen locale settings (LC_COLLATE and LC_CTYPE). If the locale is C (or equivalently POSIX), then all encodings are allowed, but for other locale settings there is only one encoding that will work properly. (On Windows, however, UTF-8 encoding can be used with any locale.) CREATE DATABASE will allow superusers to specify SQL_ASCII encoding regardless of the locale settings, but this choice is dep-
recated and may result in misbehavior of character-string functions if data that is not encoding-compatible with the locale is stored in the database.

The encoding and locale settings must match those of the template database, except when template0 is used as template. This is because other databases might contain data that does not match the specified encoding, or might contain indexes whose sort ordering is affected by LC_COLLATE and LC_CTYPE. Copying such data would result in a database that is corrupt according to the new settings. template0, however, is known to not contain any data or indexes that would be affected.

The CONNECTION LIMIT option is only enforced approximately; if two new sessions start at about the same time when just one connection “slot” remains for the database, it is possible that both will fail. Also, the limit is not enforced against superusers or background worker processes.

Examples

To create a new database:

CREATE DATABASE lusiadas;

To create a database sales owned by user salesapp with a default tablespace of salesspace:

CREATE DATABASE sales OWNER salesapp TABLESPACE salesspace;

To create a database music with a different locale:

CREATE DATABASE music
    LC_COLLATE 'sv_SE.utf8' LC_CTYPE 'sv_SE.utf8'
    TEMPLATE template0;

In this example, the TEMPLATE template0 clause is required if the specified locale is different from the one in template1. (If it is not, then specifying the locale explicitly is redundant.)

To create a database music2 with a different locale and a different character set encoding:

CREATE DATABASE music2
    LC_COLLATE 'sv_SE.iso885915' LC_CTYPE 'sv_SE.iso885915'
    ENCODING LATIN9
    TEMPLATE template0;

The specified locale and encoding settings must match, or an error will be reported.

Note that locale names are specific to the operating system, so that the above commands might not work in the same way everywhere.

Compatibility

There is no CREATE DATABASE statement in the SQL standard. Databases are equivalent to catalogs, whose creation is implementation-defined.

See Also

ALTER DATABASE, DROP DATABASE
CREATE DOMAIN

CREATE DOMAIN — define a new domain

Synopsis

CREATE DOMAIN name [ AS ] data_type
    [ COLLATE collation ]
    [ DEFAULT expression ]
    [ constraint [ ... ] ]

where constraint is:

    [ CONSTRAINT constraint_name ]
    { NOT NULL | NULL | CHECK (expression) }

Description

CREATE DOMAIN creates a new domain. A domain is essentially a data type with optional constraints (restrictions on the allowed set of values). The user who defines a domain becomes its owner.

If a schema name is given (for example, CREATE DOMAIN myschema.mydomain ...) then the domain is created in the specified schema. Otherwise it is created in the current schema. The domain name must be unique among the types and domains existing in its schema.

Domains are useful for abstracting common constraints on fields into a single location for maintenance. For example, several tables might contain email address columns, all requiring the same CHECK constraint to verify the address syntax. Define a domain rather than setting up each table's constraint individually.

To be able to create a domain, you must have USAGE privilege on the underlying type.

Parameters

name

The name (optionally schema-qualified) of a domain to be created.

data_type

The underlying data type of the domain. This can include array specifiers.

collation

An optional collation for the domain. If no collation is specified, the underlying data type's default collation is used. The underlying type must be collatable if COLLATE is specified.

DEFAULT expression

The DEFAULT clause specifies a default value for columns of the domain data type. The value is any variable-free expression (but subqueries are not allowed). The data type of the default expression must match the data type of the domain. If no default value is specified, then the default value is the null value.

The default expression will be used in any insert operation that does not specify a value for the column. If a default value is defined for a particular column, it overrides any default associated
with the domain. In turn, the domain default overrides any default value associated with the underlying data type.

CONSTRAINT constraint_name

An optional name for a constraint. If not specified, the system generates a name.

NOT NULL

Values of this domain are prevented from being null (but see notes below).

NULL

Values of this domain are allowed to be null. This is the default.

This clause is only intended for compatibility with nonstandard SQL databases. Its use is discouraged in new applications.

CHECK (expression)

CHECK clauses specify integrity constraints or tests which values of the domain must satisfy. Each constraint must be an expression producing a Boolean result. It should use the key word VALUE to refer to the value being tested. Expressions evaluating to TRUE or UNKNOWN succeed. If the expression produces a FALSE result, an error is reported and the value is not allowed to be converted to the domain type.

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than VALUE.

When a domain has multiple CHECK constraints, they will be tested in alphabetical order by name. (PostgreSQL versions before 9.5 did not honor any particular firing order for CHECK constraints.)

Notes

Domain constraints, particularly NOT NULL, are checked when converting a value to the domain type. It is possible for a column that is nominally of the domain type to read as null despite there being such a constraint. For example, this can happen in an outer-join query, if the domain column is on the nullable side of the outer join. A more subtle example is

```sql
INSERT INTO tab (domcol) VALUES ((SELECT domcol FROM tab WHERE false));
```

The empty scalar sub-SELECT will produce a null value that is considered to be of the domain type, so no further constraint checking is applied to it, and the insertion will succeed.

It is very difficult to avoid such problems, because of SQL's general assumption that a null value is a valid value of every data type. Best practice therefore is to design a domain's constraints so that a null value is allowed, and then to apply column NOT NULL constraints to columns of the domain type as needed, rather than directly to the domain type.

Examples

This example creates the us_postal_code data type and then uses the type in a table definition. A regular expression test is used to verify that the value looks like a valid US postal code:

```sql
CREATE DOMAIN us_postal_code AS TEXT
CHECK(
    VALUE ~ '^\d{5}$'
OR VALUE ~ '^\d{5}-\d{4}$'
)
CREATE TABLE us_snail_addy (  
  address_id SERIAL PRIMARY KEY,  
  street1 TEXT NOT NULL,  
  street2 TEXT,  
  street3 TEXT,  
  city TEXT NOT NULL,  
  postal us_postal_code NOT NULL
);
CREATE EVENT TRIGGER

CREATE EVENT TRIGGER — define a new event trigger

Synopsis

CREATE EVENT TRIGGER name
  ON event
  [ WHEN filter_variable IN (filter_value [, ... ]) [ AND ... ] ]
  EXECUTE { FUNCTION | PROCEDURE } function_name()

Description

CREATE EVENT TRIGGER creates a new event trigger. Whenever the designated event occurs and the WHEN condition associated with the trigger, if any, is satisfied, the trigger function will be executed. For a general introduction to event triggers, see Chapter 40. The user who creates an event trigger becomes its owner.

Parameters

name

The name to give the new trigger. This name must be unique within the database.

event

The name of the event that triggers a call to the given function. See Section 40.1 for more information on event names.

filter_variable

The name of a variable used to filter events. This makes it possible to restrict the firing of the trigger to a subset of the cases in which it is supported. Currently the only supported filter_variable is TAG.

filter_value

A list of values for the associated filter_variable for which the trigger should fire. For TAG, this means a list of command tags (e.g. ’DROP FUNCTION’).

function_name

A user-supplied function that is declared as taking no argument and returning type event_trigger.

In the syntax of CREATE EVENT TRIGGER, the keywords FUNCTION and PROCEDURE are equivalent, but the referenced function must in any case be a function, not a procedure. The use of the keyword PROCEDURE here is historical and deprecated.

Notes

Only superusers can create event triggers.

Event triggers are disabled in single-user mode (see pgsql). If an erroneous event trigger disables the database so much that you can’t even drop the trigger, restart in single-user mode and you’ll be able to do that.
CREATE EVENT TRIGGER

Examples

Forbid the execution of any DDL command:

```sql
CREATE OR REPLACE FUNCTION abort_any_command()
    RETURNS event_trigger
    LANGUAGE plpgsql
    AS $$
    BEGIN
        RAISE EXCEPTION 'command % is disabled', tg_tag;
    END;
$$;

CREATE EVENT TRIGGER abort_ddl ON ddl_command_start
    EXECUTE FUNCTION abort_any_command();
```

Compatibility

There is no `CREATE EVENT TRIGGER` statement in the SQL standard.

See Also

`ALTER EVENT TRIGGER`, `DROP EVENT TRIGGER`, `CREATE FUNCTION`
CREATE EXTENSION

CREATE EXTENSION — install an extension

Synopsis

CREATE EXTENSION [ IF NOT EXISTS ] extension_name
[ WITH ] [ SCHEMA schema_name ]
[ VERSION version ]
[ FROM old_version ]
[ CASCADE ]

Description

CREATE EXTENSION loads a new extension into the current database. There must not be an extension of the same name already loaded.

Loading an extension essentially amounts to running the extension's script file. The script will typically create new SQL objects such as functions, data types, operators and index support methods. CREATE EXTENSION additionally records the identities of all the created objects, so that they can be dropped again if DROP EXTENSION is issued.

Loading an extension requires the same privileges that would be required to create its component objects. For most extensions this means superuser or database owner privileges are needed. The user who runs CREATE EXTENSION becomes the owner of the extension for purposes of later privilege checks, as well as the owner of any objects created by the extension's script.

Parameters

IF NOT EXISTS

Do not throw an error if an extension with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing extension is anything like the one that would have been created from the currently-available script file.

extension_name

The name of the extension to be installed. PostgreSQL will create the extension using details from the file SHAREDIR/extension/extension_name.control.

schema_name

The name of the schema in which to install the extension's objects, given that the extension allows its contents to be relocated. The named schema must already exist. If not specified, and the extension's control file does not specify a schema either, the current default object creation schema is used.

If the extension specifies a schema parameter in its control file, then that schema cannot be overridden with a SCHEMA clause. Normally, an error will be raised if a SCHEMA clause is given and it conflicts with the extension's schema parameter. However, if the CASCADE clause is also given, then schema_name is ignored when it conflicts. The given schema_name will be used for installation of any needed extensions that do not specify schema in their control files.

Remember that the extension itself is not considered to be within any schema: extensions have unqualified names that must be unique database-wide. But objects belonging to the extension can be within schemas.
CREATE EXTENSION

version

The version of the extension to install. This can be written as either an identifier or a string literal. The default version is whatever is specified in the extension's control file.

old_version

FROM old_version must be specified when, and only when, you are attempting to install an extension that replaces an “old style” module that is just a collection of objects not packaged into an extension. This option causes CREATE EXTENSION to run an alternative installation script that absorbs the existing objects into the extension, instead of creating new objects. Be careful that SCHEMA specifies the schema containing these pre-existing objects.

The value to use for old_version is determined by the extension's author, and might vary if there is more than one version of the old-style module that can be upgraded into an extension. For the standard additional modules supplied with pre-9.1 PostgreSQL, use unpackaged for old_version when updating a module to extension style.

CASCADE

Automatically install any extensions that this extension depends on that are not already installed. Their dependencies are likewise automatically installed, recursively. The SCHEMA clause, if given, applies to all extensions that get installed this way. Other options of the statement are not applied to automatically-installed extensions; in particular, their default versions are always selected.

Notes

Before you can use CREATE EXTENSION to load an extension into a database, the extension's supporting files must be installed. Information about installing the extensions supplied with PostgreSQL can be found in Additional Supplied Modules.

The extensions currently available for loading can be identified from the pg_available_extensions or pg_available_extension_versions system views.

For information about writing new extensions, see Section 38.16.

Examples

Install the hstore extension into the current database:

```
CREATE EXTENSION hstore;
```

Update a pre-9.1 installation of hstore into extension style:

```
CREATE EXTENSION hstore SCHEMA public FROM unpackaged;
```

Be careful to specify the schema in which you installed the existing hstore objects.

Compatibility

CREATE EXTENSION is a PostgreSQL extension.

See Also

ALTER EXTENSION, DROP EXTENSION
CREATE FOREIGN DATA WRAPPER

CREATE FOREIGN DATA WRAPPER — define a new foreign-data wrapper

Synopsis

CREATE FOREIGN DATA WRAPPER name
   [ HANDLER handler_function | NO HANDLER ]
   [ VALIDATOR validator_function | NO VALIDATOR ]
   [ OPTIONS ( option 'value' [, ... ] ) ]

Description

CREATE FOREIGN DATA WRAPPER creates a new foreign-data wrapper. The user who defines
a foreign-data wrapper becomes its owner.

The foreign-data wrapper name must be unique within the database.

Only superusers can create foreign-data wrappers.

Parameters

name

The name of the foreign-data wrapper to be created.

HANDLER handler_function

handler_function is the name of a previously registered function that will be called to re-
trieve the execution functions for foreign tables. The handler function must take no arguments,
and its return type must be fdw_handler.

It is possible to create a foreign-data wrapper with no handler function, but foreign tables using
such a wrapper can only be declared, not accessed.

VALIDATOR validator_function

validator_function is the name of a previously registered function that will be called to
check the generic options given to the foreign-data wrapper, as well as options for foreign servers,
user mappings and foreign tables using the foreign-data wrapper. If no validator function or NO
VALIDATOR is specified, then options will not be checked at creation time. (Foreign-data wrap-
ners will possibly ignore or reject invalid option specifications at run time, depending on the im-
plementation.) The validator function must take two arguments: one of type text[], which will
contain the array of options as stored in the system catalogs, and one of type oid, which will
be the OID of the system catalog containing the options. The return type is ignored; the function
should report invalid options using the ereport (ERROR) function.

OPTIONS ( option 'value' [, ... ] )

This clause specifies options for the new foreign-data wrapper. The allowed option names and
values are specific to each foreign data wrapper and are validated using the foreign-data wrapper's
validator function. Option names must be unique.

Notes

PostgreSQL's foreign-data functionality is still under active development. Optimization of queries is
primitive (and mostly left to the wrapper, too). Thus, there is considerable room for future performance
improvements.
CREATE FOREIGN
DATA WRAPPER

Examples

Create a useless foreign-data wrapper dummy:

CREATE FOREIGN DATA WRAPPER dummy;

Create a foreign-data wrapper file with handler function file_fdw_handler:

CREATE FOREIGN DATA WRAPPER file HANDLER file_fdw_handler;

Create a foreign-data wrapper mywrapper with some options:

CREATE FOREIGN DATA WRAPPER mywrapper
  OPTIONS (debug 'true');

Compatibility

CREATE FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED), with the exception that the HANDLER and VALIDATOR clauses are extensions and the standard clauses LIBRARY and LANGUAGE are not implemented in PostgreSQL.

Note, however, that the SQL/MED functionality as a whole is not yet conforming.

See Also

ALTER FOREIGN DATA WRAPPER, DROP FOREIGN DATA WRAPPER, CREATE SERVER, CREATE USER MAPPING, CREATE FOREIGN TABLE
CREATE FOREIGN TABLE

CREATE FOREIGN TABLE — define a new foreign table

Synopsis

CREATE FOREIGN TABLE [ IF NOT EXISTS ] table_name ( [ 
   { column_name data_type [ OPTIONS ( option 'value' [, ... ] ) ] } 
[ COLLATE collation ] [ column_constraint [ ... ] ]
   | table_constraint ]
   [, ... ]
] )
[ INHERITS ( parent_table [, ... ] ) ]
SERVER server_name
[ OPTIONS ( option 'value' [, ... ] ) ]

CREATE FOREIGN TABLE [ IF NOT EXISTS ] table_name
PARTITION OF parent_table [ ( 
   { column_name [ WITH OPTIONS ] [ column_constraint [ ... ] ]
   | table_constraint ]
   [, ... ]
) ] partition_bound_spec
SERVER server_name
[ OPTIONS ( option 'value' [, ... ] ) ]

where column_constraint is:

[ CONSTRAINT constraint_name ]
{ NOT NULL | NULL |
   CHECK ( expression ) [ NO INHERIT ] |
   DEFAULT default_expr }

and table_constraint is:

[ CONSTRAINT constraint_name ]
CHECK ( expression ) [ NO INHERIT ]

Description

CREATE FOREIGN TABLE creates a new foreign table in the current database. The table will be owned by the user issuing the command.

If a schema name is given (for example, CREATE FOREIGN TABLE myschema.mytable ...) then the table is created in the specified schema. Otherwise it is created in the current schema. The name of the foreign table must be distinct from the name of any other foreign table, table, sequence, index, view, or materialized view in the same schema.

CREATE FOREIGN TABLE also automatically creates a data type that represents the composite type corresponding to one row of the foreign table. Therefore, foreign tables cannot have the same name as any existing data type in the same schema.

If PARTITION OF clause is specified then the table is created as a partition of parent_table with specified bounds.
To be able to create a foreign table, you must have USAGE privilege on the foreign server, as well as USAGE privilege on all column types used in the table.

**Parameters**

**IF NOT EXISTS**

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing relation is anything like the one that would have been created.

**table_name**

The name (optionally schema-qualified) of the table to be created.

**column_name**

The name of a column to be created in the new table.

**data_type**

The data type of the column. This can include array specifiers. For more information on the data types supported by PostgreSQL, refer to Chapter 8.

**COLLATE collation**

The COLLATE clause assigns a collation to the column (which must be of a collatable data type). If not specified, the column data type's default collation is used.

**INHERITS ( parent_table [, ... ] )**

The optional INHERITS clause specifies a list of tables from which the new foreign table automatically inherits all columns. Parent tables can be plain tables or foreign tables. See the similar form of CREATE TABLE for more details.

**CONSTRAINT constraint_name**

An optional name for a column or table constraint. If the constraint is violated, the constraint name is present in error messages, so constraint names like col must be positive can be used to communicate helpful constraint information to client applications. (Double-quotes are needed to specify constraint names that contain spaces.) If a constraint name is not specified, the system generates a name.

**NOT NULL**

The column is not allowed to contain null values.

**NULL**

The column is allowed to contain null values. This is the default.

This clause is only provided for compatibility with non-standard SQL databases. Its use is discouraged in new applications.

**CHECK ( expression ) [ NO INHERIT ]**

The CHECK clause specifies an expression producing a Boolean result which each row in the foreign table is expected to satisfy; that is, the expression should produce TRUE or UNKNOWN, never FALSE, for all rows in the foreign table. A check constraint specified as a column constraint should reference that column's value only, while an expression appearing in a table constraint can reference multiple columns.
CREATE FOREIGN TABLE

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than columns of the current row. The system column tableoid may be referenced, but not any other system column.

A constraint marked with NO INHERIT will not propagate to child tables.

DEFAULT default_expr

The DEFAULT clause assigns a default data value for the column whose column definition it appears within. The value is any variable-free expression (subqueries and cross-references to other columns in the current table are not allowed). The data type of the default expression must match the data type of the column.

The default expression will be used in any insert operation that does not specify a value for the column. If there is no default for a column, then the default is null.

server_name

The name of an existing foreign server to use for the foreign table. For details on defining a server, see CREATE SERVER.

OPTIONS ( option 'value' [, ...] )

Options to be associated with the new foreign table or one of its columns. The allowed option names and values are specific to each foreign data wrapper and are validated using the foreign-data wrapper's validator function. Duplicate option names are not allowed (although it's OK for a table option and a column option to have the same name).

Notes

Constraints on foreign tables (such as CHECK or NOT NULL clauses) are not enforced by the core PostgreSQL system, and most foreign data wrappers do not attempt to enforce them either; that is, the constraint is simply assumed to hold true. There would be little point in such enforcement since it would only apply to rows inserted or updated via the foreign table, and not to rows modified by other means, such as directly on the remote server. Instead, a constraint attached to a foreign table should represent a constraint that is being enforced by the remote server.

Some special-purpose foreign data wrappers might be the only access mechanism for the data they access, and in that case it might be appropriate for the foreign data wrapper itself to perform constraint enforcement. But you should not assume that a wrapper does that unless its documentation says so.

Although PostgreSQL does not attempt to enforce constraints on foreign tables, it does assume that they are correct for purposes of query optimization. If there are rows visible in the foreign table that do not satisfy a declared constraint, queries on the table might produce incorrect answers. It is the user's responsibility to ensure that the constraint definition matches reality.

Examples

Create foreign table films, which will be accessed through the server film_server:

```sql
CREATE FOREIGN TABLE films (   code        char(5) NOT NULL,
    title       varchar(40) NOT NULL,
    did         integer NOT NULL,
    date_prod   date,
    kind        varchar(10),
    len         interval hour to minute
)
```
CREATE FOREIGN TABLE

SERVER film_server;

Create foreign table measurement_y2016m07, which will be accessed through the server server_07, as a partition of the range partitioned table measurement:

CREATE FOREIGN TABLE measurement_y2016m07
  PARTITION OF measurement FOR VALUES FROM ('2016-07-01') TO ('2016-08-01')
  SERVER server_07;

Compatibility

The CREATE FOREIGN TABLE command largely conforms to the SQL standard; however, much as with CREATE TABLE, NULL constraints and zero-column foreign tables are permitted. The ability to specify column default values is also a PostgreSQL extension. Table inheritance, in the form defined by PostgreSQL, is nonstandard.

See Also

ALTER FOREIGN TABLE, DROP FOREIGN TABLE, CREATE TABLE, CREATE SERVER, IMPORT FOREIGN SCHEMA
CREATE FUNCTION

CREATE FUNCTION — define a new function

Synopsis

CREATE [ OR REPLACE ] FUNCTION
    name ([ [ argmode ] [ argname ] argtype [ { DEFAULT | = } default_expr ] [, ... ] ]
    [ RETURNS rettype
        | RETURNS TABLE ( column_name column_type [, ... ] ) ]
    { LANGUAGE lang_name
        | TRANSFORM { FOR TYPE type_name } [, ... ]
        | WINDOW
        | IMMUTABLE | STABLE | VOLATILE | [ NOT ] LEAKPROOF
        | CALLED ON NULL INPUT | RETURNS NULL ON NULL INPUT | STRICT
        | [ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
        | PARALLEL { UNSAFE | RESTRICTED | SAFE }
        | COST execution_cost
        | ROWS result_rows
        | SET configuration_parameter ( TO value | = value | FROM CURRENT )
        | AS 'definition'
        | AS 'obj_file', 'link_symbol'
} ...

Description

CREATE FUNCTION defines a new function. CREATE OR REPLACE FUNCTION will either create a new function, or replace an existing definition. To be able to define a function, the user must have the USAGE privilege on the language.

If a schema name is included, then the function is created in the specified schema. Otherwise it is created in the current schema. The name of the new function must not match any existing function or procedure with the same input argument types in the same schema. However, functions and procedures of different argument types can share a name (this is called overloading).

To replace the current definition of an existing function, use CREATE OR REPLACE FUNCTION. It is not possible to change the name or argument types of a function this way (if you tried, you would actually be creating a new, distinct function). Also, CREATE OR REPLACE FUNCTION will not let you change the return type of an existing function. To do that, you must drop and recreate the function. (When using OUT parameters, that means you cannot change the types of any OUT parameters except by dropping the function.)

When CREATE OR REPLACE FUNCTION is used to replace an existing function, the ownership and permissions of the function do not change. All other function properties are assigned the values specified or implied in the command. You must own the function to replace it (this includes being a member of the owning role).

If you drop and then recreate a function, the new function is not the same entity as the old; you will have to drop existing rules, views, triggers, etc. that refer to the old function. Use CREATE OR REPLACE FUNCTION to change a function definition without breaking objects that refer to the function. Also, ALTER FUNCTION can be used to change most of the auxiliary properties of an existing function.

The user that creates the function becomes the owner of the function.
CREATE FUNCTION

To be able to create a function, you must have USAGE privilege on the argument types and the return type.

**Parameters**

*name*

The name (optionally schema-qualified) of the function to create.

*argmode*

The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Only OUT arguments can follow a VARIADIC one. Also, OUT and INOUT arguments cannot be used together with the RETURNS TABLE notation.

*argname*

The name of an argument. Some languages (including SQL and PL/pgSQL) let you use the name in the function body. For other languages the name of an input argument is just extra documentation, so far as the function itself is concerned; but you can use input argument names when calling a function to improve readability (see Section 4.3). In any case, the name of an output argument is significant, because it defines the column name in the result row type. (If you omit the name for an output argument, the system will choose a default column name.)

*argtype*

The data type(s) of the function’s arguments (optionally schema-qualified), if any. The argument types can be base, composite, or domain types, or can reference the type of a table column. Depending on the implementation language it might also be allowed to specify “pseudo-types” such as cstring. Pseudo-types indicate that the actual argument type is either incompletely specified, or outside the set of ordinary SQL data types.

The type of a column is referenced by writing `table_name.column_name%TYPE`. Using this feature can sometimes help make a function independent of changes to the definition of a table.

*default_expr*

An expression to be used as default value if the parameter is not specified. The expression has to be coercible to the argument type of the parameter. Only input (including INOUT) parameters can have a default value. All input parameters following a parameter with a default value must have default values as well.

*rettype*

The return data type (optionally schema-qualified). The return type can be a base, composite, or domain type, or can reference the type of a table column. Depending on the implementation language it might also be allowed to specify “pseudo-types” such as cstring. If the function is not supposed to return a value, specify `void` as the return type.

When there are OUT or INOUT parameters, the RETURNS clause can be omitted. If present, it must agree with the result type implied by the output parameters: RECORD if there are multiple output parameters, or the same type as the single output parameter.

The SETOF modifier indicates that the function will return a set of items, rather than a single item.

The type of a column is referenced by writing `table_name.column_name%TYPE`.

*column_name*

The name of an output column in the RETURNS TABLE syntax. This is effectively another way of declaring a named OUT parameter, except that RETURNS TABLE also implies RETURNS SETOF.
**CREATE FUNCTION**

`column_type`

The data type of an output column in the `RETURNS TABLE` syntax.

`lang_name`

The name of the language that the function is implemented in. It can be `sql`, `c`, `internal`, or the name of a user-defined procedural language, e.g. `plpgsql`. Enclosing the name in single quotes is deprecated and requires matching case.

`TRANSFORM { FOR TYPE type_name } [, ... ] }`

Lists which transforms a call to the function should apply. Transforms convert between SQL types and language-specific data types; see `CREATE TRANSFORM`. Procedural language implementations usually have hardcoded knowledge of the built-in types, so those don't need to be listed here. If a procedural language implementation does not know how to handle a type and no transform is supplied, it will fall back to a default behavior for converting data types, but this depends on the implementation.

`WINDOW`  

`WINDOW` indicates that the function is a *window function* rather than a plain function. This is currently only useful for functions written in C. The `WINDOW` attribute cannot be changed when replacing an existing function definition.

`IMMUTABLE`

`STABLE`

`VOLATILE`

These attributes inform the query optimizer about the behavior of the function. At most one choice can be specified. If none of these appear, `VOLATILE` is the default assumption.

`IMMUTABLE` indicates that the function cannot modify the database and always returns the same result when given the same argument values; that is, it does not do database lookups or otherwise use information not directly present in its argument list. If this option is given, any call of the function with all-constant arguments can be immediately replaced with the function value.

`STABLE` indicates that the function cannot modify the database, and that within a single table scan it will consistently return the same result for the same argument values, but that its result could change across SQL statements. This is the appropriate selection for functions whose results depend on database lookups, parameter variables (such as the current time zone), etc. (It is inappropriate for `AFTER` triggers that wish to query rows modified by the current command.) Also note that the `current_timestamp` family of functions qualify as stable, since their values do not change within a transaction.

`VOLATILE` indicates that the function value can change even within a single table scan, so no optimizations can be made. Relatively few database functions are volatile in this sense; some examples are `random()`, `currval()`, `timeofday()`. But note that any function that has side-effects must be classified volatile, even if its result is quite predictable, to prevent calls from being optimized away; an example is `setval()`.

For additional details see Section 38.7.

`LEAKPROOF`

`LEAKPROOF` indicates that the function has no side effects. It reveals no information about its arguments other than by its return value. For example, a function which throws an error message for some argument values but not others, or which includes the argument values in any error message, is not leakproof. This affects how the system executes queries against views created with the
security_barrier option or tables with row level security enabled. The system will enforce conditions from security policies and security barrier views before any user-supplied conditions from the query itself that contain non-leakproof functions, in order to prevent the inadvertent exposure of data. Functions and operators marked as leakproof are assumed to be trustworthy, and may be executed before conditions from security policies and security barrier views. In addition, functions which do not take arguments or which are not passed any arguments from the security barrier view or table do not have to be marked as leakproof to be executed before security conditions. See CREATE VIEW and Section 41.5. This option can only be set by the superuser.

CALLED ON NULL INPUT
RETURNS NULL ON NULL INPUT
STRICT

CALLED ON NULL INPUT (the default) indicates that the function will be called normally when some of its arguments are null. It is then the function author's responsibility to check for null values if necessary and respond appropriately.

RETURNS NULL ON NULL INPUT or STRICT indicates that the function always returns null whenever any of its arguments are null. If this parameter is specified, the function is not executed when there are null arguments; instead a null result is assumed automatically.

[EXTERNAL] SECURITY INVOKER
[EXTERNAL] SECURITY DEFINER

SECURITY INVOKER indicates that the function is to be executed with the privileges of the user that calls it. That is the default. SECURITY DEFINER specifies that the function is to be executed with the privileges of the user that owns it.

The key word EXTERNAL is allowed for SQL conformance, but it is optional since, unlike in SQL, this feature applies to all functions not only external ones.

PARALLEL

PARALLEL UNSAFE indicates that the function can't be executed in parallel mode and the presence of such a function in an SQL statement forces a serial execution plan. This is the default. PARALLEL RESTRICTED indicates that the function can be executed in parallel mode, but the execution is restricted to parallel group leader. PARALLEL SAFE indicates that the function is safe to run in parallel mode without restriction.

Functions should be labeled parallel unsafe if they modify any database state, or if they make changes to the transaction such as using sub-transactions, or if they access sequences or attempt to make persistent changes to settings (e.g. setval). They should be labeled as parallel restricted if they access temporary tables, client connection state, cursors, prepared statements, or miscellaneous backend-local state which the system cannot synchronize in parallel mode (e.g. setseed cannot be executed other than by the group leader because a change made by another process would not be reflected in the leader). In general, if a function is labeled as being safe when it is restricted or unsafe, or if it is labeled as being restricted when it is in fact unsafe, it may throw errors or produce wrong answers when used in a parallel query. C-language functions could in theory exhibit totally undefined behavior if mislabeled, since there is no way for the system to protect itself against arbitrary C code, but in most likely cases the result will be no worse than for any other function. If in doubt, functions should be labeled as UNSAFE, which is the default.

COST execution_cost

A positive number giving the estimated execution cost for the function, in units of cpu_operator_cost. If the function returns a set, this is the cost per returned row. If the cost is not specified, 1 unit is assumed for C-language and internal functions, and 100 units for functions in all other languages. Larger values cause the planner to try to avoid evaluating the function more often than necessary.
CREATE FUNCTION

**ROWS** result_rows

A positive number giving the estimated number of rows that the planner should expect the function to return. This is only allowed when the function is declared to return a set. The default assumption is 1000 rows.

**configuration_parameter**

**value**

The SET clause causes the specified configuration parameter to be set to the specified value when the function is entered, and then restored to its prior value when the function exits. SET FROM CURRENT saves the value of the parameter that is current when CREATE FUNCTION is executed as the value to be applied when the function is entered.

If a SET clause is attached to a function, then the effects of a SET LOCAL command executed inside the function for the same variable are restricted to the function: the configuration parameter's prior value is still restored at function exit. However, an ordinary SET command (without LOCAL) overrides the SET clause, much as it would do for a previous SET LOCAL command: the effects of such a command will persist after function exit, unless the current transaction is rolled back.

See SET and Chapter 19 for more information about allowed parameter names and values.

**definition**

A string constant defining the function; the meaning depends on the language. It can be an internal function name, the path to an object file, an SQL command, or text in a procedural language.

It is often helpful to use dollar quoting (see Section 4.1.2.4) to write the function definition string, rather than the normal single quote syntax. Without dollar quoting, any single quotes or backslashes in the function definition must be escaped by doubling them.

**obj_file, link_symbol**

This form of the AS clause is used for dynamically loadable C language functions when the function name in the C language source code is not the same as the name of the SQL function. The string obj_file is the name of the shared library file containing the compiled C function, and is interpreted as for the LOAD command. The string link_symbol is the function's link symbol, that is, the name of the function in the C language source code. If the link symbol is omitted, it is assumed to be the same as the name of the SQL function being defined. The C names of all functions must be different, so you must give overloaded C functions different C names (for example, use the argument types as part of the C names).

When repeated CREATE FUNCTION calls refer to the same object file, the file is only loaded once per session. To unload and reload the file (perhaps during development), start a new session.

Refer to Section 38.3 for further information on writing functions.

**Overloading**

PostgreSQL allows function overloading; that is, the same name can be used for several different functions so long as they have distinct input argument types. Whether or not you use it, this capability entails security precautions when calling functions in databases where some users mistrust other users; see Section 10.3.

Two functions are considered the same if they have the same names and input argument types, ignoring any OUT parameters. Thus for example these declarations conflict:

```sql
CREATE FUNCTION foo(int) ...
```
CREATE FUNCTION foo(int, out text) ...

Functions that have different argument type lists will not be considered to conflict at creation time, but if defaults are provided they might conflict in use. For example, consider

CREATE FUNCTION foo(int) ...
CREATE FUNCTION foo(int, int default 42) ...

A call foo(10) will fail due to the ambiguity about which function should be called.

Notes

The full SQL type syntax is allowed for declaring a function's arguments and return value. However, parenthesized type modifiers (e.g., the precision field for type numeric) are discarded by CREATE FUNCTION. Thus for example CREATE FUNCTION foo (varchar(10)) ... is exactly the same as CREATE FUNCTION foo (varchar) ....

When replacing an existing function with CREATE OR REPLACE FUNCTION, there are restrictions on changing parameter names. You cannot change the name already assigned to any input parameter (although you can add names to parameters that had none before). If there is more than one output parameter, you cannot change the names of the output parameters, because that would change the column names of the anonymous composite type that describes the function's result. These restrictions are made to ensure that existing calls of the function do not stop working when it is replaced.

If a function is declared STRICT with a VARIADIC argument, the strictness check tests that the variadic array as a whole is non-null. The function will still be called if the array has null elements.

Examples

Here are some trivial examples to help you get started. For more information and examples, see Section 38.3.

CREATE FUNCTION add(integer, integer) RETURNS integer
AS 'select $1 + $2;'
LANGUAGE SQL
IMMUTABLE
RETURNS NULL ON NULL INPUT;

Increment an integer, making use of an argument name, in PL/pgSQL:

CREATE OR REPLACE FUNCTION increment(i integer) RETURNS integer AS
$$
BEGIN
  RETURN i + 1;
END;
$$ LANGUAGE plpgsql;

Return a record containing multiple output parameters:

CREATE FUNCTION dup(in int, out f1 int, out f2 text)
AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
LANGUAGE SQL;

SELECT * FROM dup(42);

You can do the same thing more verbosely with an explicitly named composite type:
CREATE TYPE dup_result AS (f1 int, f2 text);

CREATE FUNCTION dup(int) RETURNS dup_result
   AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
   LANGUAGE SQL;

SELECT * FROM dup(42);

Another way to return multiple columns is to use a TABLE function:

CREATE FUNCTION dup(int) RETURNS TABLE(f1 int, f2 text)
   AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
   LANGUAGE SQL;

SELECT * FROM dup(42);

However, a TABLE function is different from the preceding examples, because it actually returns a set of records, not just one record.

Writing SECURITY DEFINER Functions Safely

Because a SECURITY DEFINER function is executed with the privileges of the user that owns it, care is needed to ensure that the function cannot be misused. For security, search_path should be set to exclude any schemas writable by untrusted users. This prevents malicious users from creating objects (e.g., tables, functions, and operators) that mask objects intended to be used by the function. Particularly important in this regard is the temporary-table schema, which is searched first by default, and is normally writable by anyone. A secure arrangement can be obtained by forcing the temporary schema to be searched last. To do this, write pg_temp as the last entry in search_path. This function illustrates safe usage:

CREATE FUNCTION check_password(uname TEXT, pass TEXT)
RETURNS BOOLEAN AS $$
DECLARE passed BOOLEAN;
BEGIN
   SELECT (pwd = $2) INTO passed
   FROM    pwds
   WHERE   username = $1;

   RETURN passed;
END;
$$ LANGUAGE plpgsql
SECURITY DEFINER

-- Set a secure search_path: trusted schema(s), then 'pg_temp'.
SET search_path = admin, pg_temp;

This function's intention is to access a table admin.pwds. But without the SET clause, or with a SET clause mentioning only admin, the function could be subverted by creating a temporary table named pwds.

Before PostgreSQL version 8.3, the SET clause was not available, and so older functions may contain rather complicated logic to save, set, and restore search_path. The SET clause is far easier to use for this purpose.

Another point to keep in mind is that by default, execute privilege is granted to PUBLIC for newly created functions (see GRANT for more information). Frequently you will wish to restrict use of a security definer function to only some users. To do that, you must revoke the default PUBLIC
privileges and then grant execute privilege selectively. To avoid having a window where the new function is accessible to all, create it and set the privileges within a single transaction. For example:

```
BEGIN;
CREATE FUNCTION check_password(uname TEXT, pass TEXT) ... SECURITY DEFINER;
REVOKE ALL ON FUNCTION check_password(uname TEXT, pass TEXT) FROM PUBLIC;
GRANT EXECUTE ON FUNCTION check_password(uname TEXT, pass TEXT) TO admins;
COMMIT;
```

**Compatibility**

A `CREATE FUNCTION` command is defined in the SQL standard. The PostgreSQL version is similar but not fully compatible. The attributes are not portable, neither are the different available languages.

For compatibility with some other database systems, `argmode` can be written either before or after `argname`. But only the first way is standard-compliant.

For parameter defaults, the SQL standard specifies only the syntax with the `DEFAULT` key word. The syntax with `=` is used in T-SQL and Firebird.

**See Also**

`ALTER FUNCTION`, `DROP FUNCTION`, `GRANT`, `LOAD`, `REVOKE`
CREATE GROUP

CREATE GROUP — define a new database role

Synopsis

CREATE GROUP name [ [ WITH ] option [ ... ] ]

where option can be:

<table>
<thead>
<tr>
<th>SUPERUSER</th>
<th>NOSUPERUSER</th>
</tr>
</thead>
<tbody>
<tr>
<td>CREATEDB</td>
<td>NOCREATEDB</td>
</tr>
<tr>
<td>CREATEROLE</td>
<td>NOCREATEROLE</td>
</tr>
<tr>
<td>INHERIT</td>
<td>NOINHERIT</td>
</tr>
<tr>
<td>LOGIN</td>
<td>NOLOGIN</td>
</tr>
<tr>
<td>[ ENCRYPTED ] PASSWORD 'password'</td>
<td></td>
</tr>
<tr>
<td>VALID UNTIL 'timestamp'</td>
<td></td>
</tr>
<tr>
<td>IN ROLE role_name [, ...]</td>
<td></td>
</tr>
<tr>
<td>IN GROUP role_name [, ...]</td>
<td></td>
</tr>
<tr>
<td>ROLE role_name [, ...]</td>
<td></td>
</tr>
<tr>
<td>ADMIN role_name [, ...]</td>
<td></td>
</tr>
<tr>
<td>USER role_name [, ...]</td>
<td></td>
</tr>
<tr>
<td>SYSID uid</td>
<td></td>
</tr>
</tbody>
</table>

Description

CREATE GROUP is now an alias for CREATE ROLE.

Compatibility

There is no CREATE GROUP statement in the SQL standard.

See Also

CREATE ROLE
CREATE INDEX

CREATE INDEX — define a new index

Synopsis

CREATE [ UNIQUE ] INDEX [ CONCURRENTLY ] [ [ IF NOT EXISTS ] name ]
ON [ ONLY ] table_name [ USING method ]
( { column_name | ( expression ) } [ COLLATE collation ]
[ opclass ] [ ASC | DESC ] [ NULLS { FIRST | LAST } ] [, ... ] )
[ INCLUDE ( column_name [, ... ] ) ]
[ WITH ( storage_parameter = value [, ... ] ) ]
[ TABLESPACE tablespace_name ]
[ WHERE predicate ]

Description

CREATE INDEX constructs an index on the specified column(s) of the specified relation, which can be a table or a materialized view. Indexes are primarily used to enhance database performance (though inappropriate use can result in slower performance).

The key field(s) for the index are specified as column names, or alternatively as expressions written in parentheses. Multiple fields can be specified if the index method supports multicolumn indexes.

An index field can be an expression computed from the values of one or more columns of the table row. This feature can be used to obtain fast access to data based on some transformation of the basic data. For example, an index computed on upper(col) would allow the clause WHERE upper(col) = 'JIM' to use an index.

PostgreSQL provides the index methods B-tree, hash, GiST, SP-GiST, GIN, and BRIN. Users can also define their own index methods, but that is fairly complicated.

When the WHERE clause is present, a partial index is created. A partial index is an index that contains entries for only a portion of a table, usually a portion that is more useful for indexing than the rest of the table. For example, if you have a table that contains both billed and unbilled orders where the unbilled orders take up a small fraction of the total table and yet that is an often used section, you can improve performance by creating an index on just that portion. Another possible application is to use WHERE with UNIQUE to enforce uniqueness over a subset of a table. See Section 11.8 for more discussion.

The expression used in the WHERE clause can refer only to columns of the underlying table, but it can use all columns, not just the ones being indexed. Presently, subqueries and aggregate expressions are also forbidden in WHERE. The same restrictions apply to index fields that are expressions.

All functions and operators used in an index definition must be “immutable”, that is, their results must depend only on their arguments and never on any outside influence (such as the contents of another table or the current time). This restriction ensures that the behavior of the index is well-defined. To use a user-defined function in an index expression or WHERE clause, remember to mark the function immutable when you create it.

Parameters

UNIQUE

Causes the system to check for duplicate values in the table when the index is created (if data already exist) and each time data is added. Attempts to insert or update data which would result in duplicate entries will generate an error.
CREATE INDEX

Additional restrictions apply when unique indexes are applied to partitioned tables; see CREATE TABLE.

CONCURRENTLY

When this option is used, PostgreSQL will build the index without taking any locks that prevent concurrent inserts, updates, or deletes on the table; whereas a standard index build locks out writes (but not reads) on the table until it's done. There are several caveats to be aware of when using this option — see Building Indexes Concurrently.

IF NOT EXISTS

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing index is anything like the one that would have been created. Index name is required when IF NOT EXISTS is specified.

INCLUDE

The optional INCLUDE clause specifies a list of columns which will be included in the index as non-key columns. A non-key column cannot be used in an index scan search qualification, and it is disregarded for purposes of any uniqueness or exclusion constraint enforced by the index. However, an index-only scan can return the contents of non-key columns without having to visit the index's table, since they are available directly from the index entry. Thus, addition of non-key columns allows index-only scans to be used for queries that otherwise could not use them.

It's wise to be conservative about adding non-key columns to an index, especially wide columns. If an index tuple exceeds the maximum size allowed for the index type, data insertion will fail. In any case, non-key columns duplicate data from the index's table and bloat the size of the index, thus potentially slowing searches.

Columns listed in the INCLUDE clause don't need appropriate operator classes; the clause can include columns whose data types don't have operator classes defined for a given access method.

Expressions are not supported as included columns since they cannot be used in index-only scans.

Currently, only the B-tree index access method supports this feature. In B-tree indexes, the values of columns listed in the INCLUDE clause are included in leaf tuples which correspond to heap tuples, but are not included in upper-level index entries used for tree navigation.

name

The name of the index to be created. No schema name can be included here; the index is always created in the same schema as its parent table. If the name is omitted, PostgreSQL chooses a suitable name based on the parent table's name and the indexed column name(s).

ONLY

Indicates not to recurse creating indexes on partitions, if the table is partitioned. The default is to recurse.

table_name

The name (possibly schema-qualified) of the table to be indexed.

method

The name of the index method to be used. Choices are btree, hash, gist, spgist, gin, and brin. The default method is btree.

column_name

The name of a column of the table.
CREATE INDEX

expression

An expression based on one or more columns of the table. The expression usually must be written with surrounding parentheses, as shown in the syntax. However, the parentheses can be omitted if the expression has the form of a function call.

collation

The name of the collation to use for the index. By default, the index uses the collation declared for the column to be indexed or the result collation of the expression to be indexed. Indexes with non-default collations can be useful for queries that involve expressions using non-default collations.

opclass

The name of an operator class. See below for details.

ASC

Specifies ascending sort order (which is the default).

DESC

Specifies descending sort order.

NULLS FIRST

Specifies that nulls sort before non-nulls. This is the default when DESC is specified.

NULLS LAST

Specifies that nulls sort after non-nulls. This is the default when DESC is not specified.

storage_parameter

The name of an index-method-specific storage parameter. See Index Storage Parameters for details.

tablespace_name

The tablespace in which to create the index. If not specified, default_tablespace is consulted, or temp_tablespaces for indexes on temporary tables.

predicate

The constraint expression for a partial index.

Index Storage Parameters

The optional WITH clause specifies storage parameters for the index. Each index method has its own set of allowed storage parameters. The B-tree, hash, GiST and SP-GiST index methods all accept this parameter:

fillfactor

The fillfactor for an index is a percentage that determines how full the index method will try to pack index pages. For B-trees, leaf pages are filled to this percentage during initial index build, and also when extending the index at the right (adding new largest key values). If pages subsequently become completely full, they will be split, leading to gradual degradation in the index's efficiency. B-trees use a default fillfactor of 90, but any integer value from 10 to 100 can be selected. If the table is static then fillfactor 100 is best to minimize the index's physical size, but for heavily updated tables a smaller fillfactor is better to minimize the need for page splits. The other index methods use fillfactor in different but roughly analogous ways; the default fillfactor varies between methods.
CREATE INDEX

B-tree indexes additionally accept this parameter:

vacuum_cleanup_index_scale_factor

Per-index value for vacuum_cleanup_index_scale_factor.

GiST indexes additionally accept this parameter:

buffering

Determines whether the buffering build technique described in Section 64.4.1 is used to build the index. With OFF it is disabled, with ON it is enabled, and with AUTO it is initially disabled, but turned on on-the-fly once the index size reaches effective_cache_size. The default is AUTO.

GIN indexes accept different parameters:

fastupdate

This setting controls usage of the fast update technique described in Section 66.4.1. It is a Boolean parameter: ON enables fast update, OFF disables it. (Alternative spellings of ON and OFF are allowed as described in Section 19.1.) The default is ON.

Note

Turning fastupdate off via ALTER INDEX prevents future insertions from going into the list of pending index entries, but does not in itself flush previous entries. You might want to VACUUM the table or call gin_clean_pending_list function afterward to ensure the pending list is emptied.

gin_pending_list_limit

Custom gin_pending_list_limit parameter. This value is specified in kilobytes.

BRIN indexes accept different parameters:

pages_per_range

Defines the number of table blocks that make up one block range for each entry of a BRIN index (see Section 67.1 for more details). The default is 128.

autosummarize

Defines whether a summarization run is invoked for the previous page range whenever an insertion is detected on the next one.

Building Indexes Concurrently

Creating an index can interfere with regular operation of a database. Normally PostgreSQL locks the table to be indexed against writes and performs the entire index build with a single scan of the table. Other transactions can still read the table, but if they try to insert, update, or delete rows in the table they will block until the index build is finished. This could have a severe effect if the system is a live production database. Very large tables can take many hours to be indexed, and even for smaller tables, an index build can lock out writers for periods that are unacceptably long for a production system.

PostgreSQL supports building indexes without locking out writes. This method is invoked by specifying the CONCURRENTLY option of CREATE_INDEX. When this option is used, PostgreSQL must perform two scans of the table, and in addition it must wait for all existing transactions that could potentially modify or use the index to terminate. Thus this method requires more total work than a standard index build and takes significantly longer to complete. However, since it allows normal op-
erations to continue while the index is built, this method is useful for adding new indexes in a production environment. Of course, the extra CPU and I/O load imposed by the index creation might slow other operations.

In a concurrent index build, the index is actually entered into the system catalogs in one transaction, then two table scans occur in two more transactions. Before each table scan, the index build must wait for existing transactions that have modified the table to terminate. After the second scan, the index build must wait for any transactions that have a snapshot (see Chapter 13) predating the second scan to terminate. Then finally the index can be marked ready for use, and the CREATE INDEX command terminates. Even then, however, the index may not be immediately usable for queries: in the worst case, it cannot be used as long as transactions exist that predate the start of the index build.

If a problem arises while scanning the table, such as a deadlock or a uniqueness violation in a unique index, the CREATE_INDEX command will fail but leave behind an “invalid” index. This index will be ignored for querying purposes because it might be incomplete; however it will still consume update overhead. The psql \d command will report such an index as INVALID:

```
postgres=# \d tab
Table "public.tab"
 Column | Type    | Collation | Nullable | Default
--------+---------+-----------+----------+---------
  col    | integer |           |          |
Indexes:
   "idx" btree (col) INVALID
```

The recommended recovery method in such cases is to drop the index and try again to perform CREATE_INDEX CONCURRENTLY. (Another possibility is to rebuild the index with REINDEX. However, since REINDEX does not support concurrent builds, this option is unlikely to seem attractive.)

Another caveat when building a unique index concurrently is that the uniqueness constraint is already being enforced against other transactions when the second table scan begins. This means that constraint violations could be reported in other queries prior to the index becoming available for use, or even in cases where the index build eventually fails. Also, if a failure does occur in the second scan, the “invalid” index continues to enforce its uniqueness constraint afterwards.

Concurrent builds of expression indexes and partial indexes are supported. Errors occurring in the evaluation of these expressions could cause behavior similar to that described above for unique constraint violations.

Regular index builds permit other regular index builds on the same table to occur simultaneously, but only one concurrent index build can occur on a table at a time. In either case, schema modification of the table is not allowed while the index is being built. Another difference is that a regular CREATE_INDEX command can be performed within a transaction block, but CREATE_INDEX CONCURRENTLY cannot.

Notes

See Chapter 11 for information about when indexes can be used, when they are not used, and in which particular situations they can be useful.

Currently, only the B-tree, GiST, GIN, and BRIN index methods support multicolumn indexes. Up to 32 fields can be specified by default. (This limit can be altered when building PostgreSQL.) Only B-tree currently supports unique indexes.

An operator class can be specified for each column of an index. The operator class identifies the operators to be used by the index for that column. For example, a B-tree index on four-byte integers would use the int4_ops class; this operator class includes comparison functions for four-byte integers. In practice the default operator class for the column’s data type is usually sufficient. The main
CREATE INDEX

point of having operator classes is that for some data types, there could be more than one meaningful ordering. For example, we might want to sort a complex-number data type either by absolute value or by real part. We could do this by defining two operator classes for the data type and then selecting the proper class when creating an index. More information about operator classes is in Section 11.9 and in Section 38.15.

When `CREATE INDEX` is invoked on a partitioned table, the default behavior is to recurse to all partitions to ensure they all have matching indexes. Each partition is first checked to determine whether an equivalent index already exists, and if so, that index will become attached as a partition index to the index being created, which will become its parent index. If no matching index exists, a new index will be created and automatically attached; the name of the new index in each partition will be determined as if no index name had been specified in the command. If the `ONLY` option is specified, no recursion is done, and the index is marked invalid. (`ALTER INDEX ... ATTACH PARTITION` marks the index valid, once all partitions acquire matching indexes.) Note, however, that any partition that is created in the future using `CREATE TABLE ... PARTITION OF` will automatically have a matching index, regardless of whether `ONLY` is specified.

For index methods that support ordered scans (currently, only B-tree), the optional clauses `ASC`, `DESC`, `NULLS FIRST`, and/or `NULLS LAST` can be specified to modify the sort ordering of the index. Since an ordered index can be scanned either forward or backward, it is not normally useful to create a single-column `DESC` index — that sort ordering is already available with a regular index. The value of these options is that multicolumn indexes can be created that match the sort ordering requested by a mixed-ordering query, such as `SELECT ... ORDER BY x ASC, y DESC`. The `NULLS` options are useful if you need to support “nulls sort low” behavior, rather than the default “nulls sort high”, in queries that depend on indexes to avoid sorting steps.

For most index methods, the speed of creating an index is dependent on the setting of maintenance_work_mem. Larger values will reduce the time needed for index creation, so long as you don't make it larger than the amount of memory really available, which would drive the machine into swapping.

PostgreSQL can build indexes while leveraging multiple CPUs in order to process the table rows faster. This feature is known as parallel index build. For index methods that support building indexes in parallel (currently, only B-tree), maintenance_work_mem specifies the maximum amount of memory that can be used by each index build operation as a whole, regardless of how many worker processes were started. Generally, a cost model automatically determines how many worker processes should be requested, if any.

Parallel index builds may benefit from increasing maintenance_work_mem where an equivalent serial index build will see little or no benefit. Note that maintenance_work_mem may influence the number of worker processes requested, since parallel workers must have at least a 32MB share of the total maintenance_work_mem budget. There must also be a remaining 32MB share for the leader process. Increasing max_parallel_maintenance_workers may allow more workers to be used, which will reduce the time needed for index creation, so long as the index build is not already I/O bound. Of course, there should also be sufficient CPU capacity that would otherwise lie idle.

Setting a value for parallel_workers via `ALTER TABLE` directly controls how many parallel worker processes will be requested by a `CREATE INDEX` against the table. This bypasses the cost model completely, and prevents maintenance_work_mem from affecting how many parallel workers are requested. Setting `parallel_workers` to 0 via `ALTER TABLE` will disable parallel index builds on the table in all cases.

**Tip**

You might want to reset `parallel_workers` after setting it as part of tuning an index build. This avoids inadvertent changes to query plans, since `parallel_workers` affects all parallel table scans.
While `CREATE INDEX` with the `CONCURRENTLY` option supports parallel builds without special restrictions, only the first table scan is actually performed in parallel.

Use `DROP INDEX` to remove an index.

Prior releases of PostgreSQL also had an R-tree index method. This method has been removed because it had no significant advantages over the GiST method. If `USING rtree` is specified, `CREATE INDEX` will interpret it as `USING gist`, to simplify conversion of old databases to GiST.

**Examples**

To create a unique B-tree index on the column `title` in the table `films`:

```sql
CREATE UNIQUE INDEX title_idx ON films (title);
```

To create a unique B-tree index on the column `title` with included columns `director` and `rating` in the table `films`:

```sql
CREATE UNIQUE INDEX title_idx ON films (title) INCLUDE (director, rating);
```

To create an index on the expression `lower(title)`, allowing efficient case-insensitive searches:

```sql
CREATE INDEX ON films (lower(title));
```

(In this example we have chosen to omit the index name, so the system will choose a name, typically `films_lower_idx`.)

To create an index with non-default collation:

```sql
CREATE INDEX title_idx_german ON films (title COLLATE "de_DE");
```

To create an index with non-default sort ordering of nulls:

```sql
CREATE INDEX title_idx_nulls_low ON films (title NULLS FIRST);
```

To create an index with non-default fill factor:

```sql
CREATE UNIQUE INDEX title_idx ON films (title) WITH (fillfactor = 70);
```

To create a GIN index with fast updates disabled:

```sql
CREATE INDEX gin_idx ON documents_table USING GIN (locations) WITH (fastupdate = off);
```

To create an index on the column `code` in the table `films` and have the index reside in the tablespace `indexspace`:

```sql
CREATE INDEX code_idx ON films (code) TABLESPACE indexspace;
```

To create a GiST index on a point attribute so that we can efficiently use box operators on the result of the conversion function:
CREATE INDEX pointloc
    ON points USING gist (box(location,location));
SELECT * FROM points
    WHERE box(location,location) && '(0,0),(1,1)::box;

To create an index without locking out writes to the table:

CREATE INDEX CONCURRENTLY sales_quantity_index ON sales_table
    (quantity);

Compatibility

CREATE_INDEX is a PostgreSQL language extension. There are no provisions for indexes in the SQL standard.

See Also

ALTER INDEX, DROP INDEX
CREATE LANGUAGE

CREATE LANGUAGE — define a new procedural language

Synopsis

CREATE [ OR REPLACE ] [ PROCEDURAL ] LANGUAGE name
CREATE [ OR REPLACE ] [ TRUSTED ] [ PROCEDURAL ] LANGUAGE name
    HANDLER call_handler [ INLINE inline_handler ]
    [ VALIDATOR valfunction ]

Description

CREATE LANGUAGE registers a new procedural language with a PostgreSQL database. Subsequently, functions and procedures can be defined in this new language.

Note

As of PostgreSQL 9.1, most procedural languages have been made into “extensions”, and should therefore be installed with CREATE EXTENSION not CREATE LANGUAGE. Direct use of CREATE LANGUAGE should now be confined to extension installation scripts. If you have a “bare” language in your database, perhaps as a result of an upgrade, you can convert it to an extension using CREATE EXTENSION language FROM unpackaged.

CREATE LANGUAGE effectively associates the language name with handler function(s) that are responsible for executing functions written in the language. Refer to Chapter 56 for more information about language handlers.

There are two forms of the CREATE LANGUAGE command. In the first form, the user supplies just the name of the desired language, and the PostgreSQL server consults the pg_pltemplate system catalog to determine the correct parameters. In the second form, the user supplies the language parameters along with the language name. The second form can be used to create a language that is not defined in pg_pltemplate, but this approach is considered obsolescent.

When the server finds an entry in the pg_pltemplate catalog for the given language name, it will use the catalog data even if the command includes language parameters. This behavior simplifies loading of old dump files, which are likely to contain out-of-date information about language support functions.

Ordinarily, the user must have the PostgreSQL superuser privilege to register a new language. However, the owner of a database can register a new language within that database if the language is listed in the pg_pltemplate catalog and is marked as allowed to be created by database owners (tm_pldbacreate is true). The default is that trusted languages can be created by database owners, but this can be adjusted by superusers by modifying the contents of pg_pltemplate. The creator of a language becomes its owner and can later drop it, rename it, or assign it to a new owner.

CREATE OR REPLACE LANGUAGE will either create a new language, or replace an existing definition. If the language already exists, its parameters are updated according to the values specified or taken from pg_pltemplate, but the language’s ownership and permissions settings do not change, and any existing functions written in the language are assumed to still be valid. In addition to the normal privilege requirements for creating a language, the user must be superuser or owner of the existing language. The REPLACE case is mainly meant to be used to ensure that the language exists. If the
CREATE LANGUAGE

language has a pg_pltemplate entry then REPLACE will not actually change anything about an existing definition, except in the unusual case where the pg_pltemplate entry has been modified since the language was created.

Parameters

TRUSTED

TRUSTED specifies that the language does not grant access to data that the user would not otherwise have. If this key word is omitted when registering the language, only users with the PostgreSQL superuser privilege can use this language to create new functions.

PROCEDURAL

This is a noise word.

name

The name of the new procedural language. The name must be unique among the languages in the database.

For backward compatibility, the name can be enclosed by single quotes.

HANDLER call_handler

call_handler is the name of a previously registered function that will be called to execute the procedural language's functions. The call handler for a procedural language must be written in a compiled language such as C with version 1 call convention and registered with PostgreSQL as a function taking no arguments and returning the language_handler type, a placeholder type that is simply used to identify the function as a call handler.

INLINE inline_handler

inline_handler is the name of a previously registered function that will be called to execute an anonymous code block (DO command) in this language. If no inline_handler function is specified, the language does not support anonymous code blocks. The handler function must take one argument of type internal, which will be the DO command's internal representation, and it will typically return void. The return value of the handler is ignored.

VALIDATOR valfunction

valfunction is the name of a previously registered function that will be called when a new function in the language is created, to validate the new function. If no validator function is specified, then a new function will not be checked when it is created. The validator function must take one argument of type oid, which will be the OID of the to-be-created function, and will typically return void. 

A validator function would typically inspect the function body for syntactical correctness, but it can also look at other properties of the function, for example if the language cannot handle certain argument types. To signal an error, the validator function should use the ereport() function. The return value of the function is ignored.

The TRUSTED option and the support function name(s) are ignored if the server has an entry for the specified language name in pg_pltemplate.

Notes

Use DROP LANGUAGE to drop procedural languages.

The system catalog pg_language (see Section 52.29) records information about the currently installed languages. Also, the psql command \dL lists the installed languages.
To create functions in a procedural language, a user must have the USAGE privilege for the language. By default, USAGE is granted to PUBLIC (i.e., everyone) for trusted languages. This can be revoked if desired.

Procedural languages are local to individual databases. However, a language can be installed into the template1 database, which will cause it to be available automatically in all subsequently-created databases.

The call handler function, the inline handler function (if any), and the validator function (if any) must already exist if the server does not have an entry for the language in pg_pltemplate. But when there is an entry, the functions need not already exist; they will be automatically defined if not present in the database. (This might result in CREATE LANGUAGE failing, if the shared library that implements the language is not available in the installation.)

In PostgreSQL versions before 7.3, it was necessary to declare handler functions as returning the placeholder type opaque, rather than language_handler. To support loading of old dump files, CREATE LANGUAGE will accept a function declared as returning opaque, but it will issue a notice and change the function's declared return type to language_handler.

**Examples**

The preferred way of creating any of the standard procedural languages is just:

```
CREATE LANGUAGE plperl;
```

For a language not known in the pg_pltemplate catalog, a sequence such as this is needed:

```
CREATE FUNCTION plsample_call_handler() RETURNS language_handler
AS '$libdir/plsample'
LANGUAGE C;
CREATE LANGUAGE plsample
HANDLER plsample_call_handler;
```

**Compatibility**

CREATE LANGUAGE is a PostgreSQL extension.

**See Also**

ALTER LANGUAGE, CREATE FUNCTION, DROP LANGUAGE, GRANT, REVOKE
CREATE MATERIALIZED VIEW

CREATE MATERIALIZED VIEW — define a new materialized view

Synopsis

CREATE MATERIALIZED VIEW [ IF NOT EXISTS ] table_name
  [ (column_name [, ...] ) ]
  [ WITH ( storage_parameter [= value] [, ... ] ) ]
  [ TABLESPACE tablespace_name ]
AS query
  [ WITH [ NO ] DATA ]

Description

CREATE MATERIALIZED VIEW defines a materialized view of a query. The query is executed and
used to populate the view at the time the command is issued (unless WITH NO DATA is used) and
may be refreshed later using REFRESH MATERIALIZED VIEW.

CREATE MATERIALIZED VIEW is similar to CREATE TABLE AS, except that it also remembers
the query used to initialize the view, so that it can be refreshed later upon demand. A materialized
view has many of the same properties as a table, but there is no support for temporary materialized
views or automatic generation of OIDs.

Parameters

IF NOT EXISTS

Do not throw an error if a materialized view with the same name already exists. A notice is issued
in this case. Note that there is no guarantee that the existing materialized view is anything like
the one that would have been created.

table_name

The name (optionally schema-qualified) of the materialized view to be created.

column_name

The name of a column in the new materialized view. If column names are not provided, they are
taken from the output column names of the query.

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for the new materialized view; see Storage Pa-
rameters for more information. All parameters supported for CREATE TABLE are also supported
for CREATE MATERIALIZED VIEW with the exception of OIDS. See CREATE TABLE for
more information.

TABLESPACE tablespace_name

The tablespace_name is the name of the tablespace in which the new materialized view is
to be created. If not specified, default_tablespace is consulted.

query

A SELECT, TABLE, or VALUES command. This query will run within a security-restricted
operation; in particular, calls to functions that themselves create temporary tables will fail.
WITH [ NO ] DATA

This clause specifies whether or not the materialized view should be populated at creation time. If not, the materialized view will be flagged as unscannable and cannot be queried until REFRESH MATERIALIZED VIEW is used.

Compatibility

CREATE MATERIALIZED VIEW is a PostgreSQL extension.

See Also

ALTER MATERIALIZED VIEW, CREATE TABLE AS, CREATE VIEW, DROP MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
CREATE OPERATOR

CREATE OPERATOR — define a new operator

Synopsis

CREATE OPERATOR name {
   (FUNCTION|PROCEDURE) = function_name
   [, LEFTARG = left_type ] [, RIGHTARG = right_type ]
   [, COMMUTATOR = com_op ] [, NEGATOR = neg_op ]
   [, RESTRICT = res_proc ] [, JOIN = join_proc ]
   [, HASHES ] [, MERGES ]
}

Description

CREATE OPERATOR defines a new operator, name. The user who defines an operator becomes its owner. If a schema name is given then the operator is created in the specified schema. Otherwise it is created in the current schema.

The operator name is a sequence of up to NAMEDATALEN-1 (63 by default) characters from the following list:

+ - */ < > = ~ ! @ # % ^ & | ` ?

There are a few restrictions on your choice of name:

• `--` and `/*` cannot appear anywhere in an operator name, since they will be taken as the start of a comment.

• A multicharacter operator name cannot end in `+` or `-`, unless the name also contains at least one of these characters:

```
~ ! @ # % ^ & | ` ?
```

For example, `@--` is an allowed operator name, but `*--` is not. This restriction allows PostgreSQL to parse SQL-compliant commands without requiring spaces between tokens.

• The use of `=>` as an operator name is deprecated. It may be disallowed altogether in a future release.

The operator `!=` is mapped to `<>` on input, so these two names are always equivalent.

At least one of LEFTARG and RIGHTARG must be defined. For binary operators, both must be defined. For right unary operators, only LEFTARG should be defined, while for left unary operators only RIGHTARG should be defined.

The function_name function must have been previously defined using CREATE FUNCTION and must be defined to accept the correct number of arguments (either one or two) of the indicated types.

In the syntax of CREATE OPERATOR, the keywords FUNCTION and PROCEDURE are equivalent, but the referenced function must in any case be a function, not a procedure. The use of the keyword PROCEDURE here is historical and deprecated.

The other clauses specify optional operator optimization clauses. Their meaning is detailed in Section 38.14.
To be able to create an operator, you must have USAGE privilege on the argument types and the return type, as well as EXECUTE privilege on the underlying function. If a commutator or negator operator is specified, you must own these operators.

Parameters

name
The name of the operator to be defined. See above for allowable characters. The name can be schema-qualified, for example CREATE OPERATOR myschema.+ (...). If not, then the operator is created in the current schema. Two operators in the same schema can have the same name if they operate on different data types. This is called overloading.

function_name
The function used to implement this operator.

left_type
The data type of the operator's left operand, if any. This option would be omitted for a left-unary operator.

right_type
The data type of the operator's right operand, if any. This option would be omitted for a right-unary operator.

com_op
The commutator of this operator.

neg_op
The negator of this operator.

res_proc
The restriction selectivity estimator function for this operator.

join_proc
The join selectivity estimator function for this operator.

HASHES
Indicates this operator can support a hash join.

MERGES
Indicates this operator can support a merge join.

To give a schema-qualified operator name in com_op or the other optional arguments, use the OPERATOR() syntax, for example:

COMMUTATOR = OPERATOR(myschema.===),

Notes

Refer to Section 38.13 for further information.
It is not possible to specify an operator's lexical precedence in `CREATE OPERATOR`, because the parser's precedence behavior is hard-wired. See Section 4.1.6 for precedence details.

The obsolete options `SORT1`, `SORT2`, `LT_CMP`, and `GT_CMP` were formerly used to specify the names of sort operators associated with a merge-joinable operator. This is no longer necessary, since information about associated operators is found by looking at B-tree operator families instead. If one of these options is given, it is ignored except for implicitly setting `MERGES` true.

Use `DROP OPERATOR` to delete user-defined operators from a database. Use `ALTER OPERATOR` to modify operators in a database.

### Examples

The following command defines a new operator, `area-equality`, for the data type `box`:

```sql
CREATE OPERATOR === (
    LEFTARG = box,
    RIGHTARG = box,
    FUNCTION = area_equal_function,
    COMMUTATOR = ===,
    NEGATOR = !==,
    RESTRICT = area_restriction_function,
    JOIN = area_join_function,
    MERGES
)
```

### Compatibility

`CREATE OPERATOR` is a PostgreSQL extension. There are no provisions for user-defined operators in the SQL standard.

### See Also

`ALTER OPERATOR`, `CREATE OPERATOR CLASS`, `DROP OPERATOR`
CREATE OPERATOR CLASS

CREATE OPERATOR CLASS — define a new operator class

Synopsis

CREATE OPERATOR CLASS name [ DEFAULT ] FOR TYPE data_type
  USING index_method [ FAMILY family_name ] AS
  { OPERATOR strategy_number operator_name [ ( op_type, op_type ]
    ] [ FOR SEARCH [ FOR ORDER BY sort_family_name ]
      | FUNCTION support_number [ ( op_type [ , op_type ]
    ] } ) function_name ( argument_type [, ... ] )
  | STORAGE storage_type
} [ , ... ]

Description

CREATE OPERATOR CLASS creates a new operator class. An operator class defines how a particular data type can be used with an index. The operator class specifies that certain operators will fill particular roles or “strategies” for this data type and this index method. The operator class also specifies the support functions to be used by the index method when the operator class is selected for an index column. All the operators and functions used by an operator class must be defined before the operator class can be created.

If a schema name is given then the operator class is created in the specified schema. Otherwise it is created in the current schema. Two operator classes in the same schema can have the same name only if they are for different index methods.

The user who defines an operator class becomes its owner. Presently, the creating user must be a superuser. (This restriction is made because an erroneous operator class definition could confuse or even crash the server.)

CREATE OPERATOR CLASS does not presently check whether the operator class definition includes all the operators and functions required by the index method, nor whether the operators and functions form a self-consistent set. It is the user's responsibility to define a valid operator class.

Related operator classes can be grouped into operator families. To add a new operator class to an existing family, specify the FAMILY option in CREATE OPERATOR CLASS. Without this option, the new class is placed into a family named the same as the new class (creating that family if it doesn't already exist).

Refer to Section 38.15 for further information.

Parameters

name

The name of the operator class to be created. The name can be schema-qualified.

DEFAULT

If present, the operator class will become the default operator class for its data type. At most one operator class can be the default for a specific data type and index method.

data_type

The column data type that this operator class is for.
CREATE OPERATOR CLASS

index_method

The name of the index method this operator class is for.

family_name

The name of the existing operator family to add this operator class to. If not specified, a family named the same as the operator class is used (creating it, if it doesn't already exist).

strategy_number

The index method's strategy number for an operator associated with the operator class.

operator_name

The name (optionally schema-qualified) of an operator associated with the operator class.

op_type

In an OPERATOR clause, the operand data type(s) of the operator, or NONE to signify a left-unary or right-unary operator. The operand data types can be omitted in the normal case where they are the same as the operator class's data type.

In a FUNCTION clause, the operand data type(s) the function is intended to support, if different from the input data type(s) of the function (for B-tree comparison functions and hash functions) or the class's data type (for B-tree sort support functions and all functions in GiST, SP-GiST, GIN and BRIN operator classes). These defaults are correct, and so op_type need not be specified in FUNCTION clauses, except for the case of a B-tree sort support function that is meant to support cross-data-type comparisons.

sort_family_name

The name (optionally schema-qualified) of an existing btree operator family that describes the sort ordering associated with an ordering operator.

If neither FOR SEARCH nor FOR ORDER BY is specified, FOR SEARCH is the default.

support_number

The index method's support function number for a function associated with the operator class.

function_name

The name (optionally schema-qualified) of a function that is an index method support function for the operator class.

argument_type

The parameter data type(s) of the function.

storage_type

The data type actually stored in the index. Normally this is the same as the column data type, but some index methods (currently GiST, GIN and BRIN) allow it to be different. The STORAGE clause must be omitted unless the index method allows a different type to be used. If the column data_type is specified as anyarray, the storage_type can be declared as anyelement to indicate that the index entries are members of the element type belonging to the actual array type that each particular index is created for.

The OPERATOR, FUNCTION, and STORAGE clauses can appear in any order.
Notes

Because the index machinery does not check access permissions on functions before using them, including a function or operator in an operator class is tantamount to granting public execute permission on it. This is usually not an issue for the sorts of functions that are useful in an operator class.

The operators should not be defined by SQL functions. A SQL function is likely to be inlined into the calling query, which will prevent the optimizer from recognizing that the query matches an index.

Before PostgreSQL 8.4, the OPERATOR clause could include a RECHECK option. This is no longer supported because whether an index operator is “lossy” is now determined on-the-fly at run time. This allows efficient handling of cases where an operator might or might not be lossy.

Examples

The following example command defines a GiST index operator class for the data type _int4 (array of int4). See the intarray module for the complete example.

```
CREATE OPERATOR CLASS gist__int_ops
  DEFAULT FOR TYPE _int4 USING gist AS
  OPERATOR 3 &&, 
  OPERATOR 6 = (anyarray, anyarray),
  OPERATOR 7 >,
  OPERATOR 8 <@,
  OPERATOR 20 @@ (_int4, query_int),
  FUNCTION 1 g_int_consistent (internal, _int4, smallint, oid, internal),
  FUNCTION 2 g_int_union (internal, internal),
  FUNCTION 3 g_int_compress (internal),
  FUNCTION 4 g_int_decompress (internal),
  FUNCTION 5 g_int_penalty (internal, internal, internal),
  FUNCTION 6 g_int_picksplit (internal, internal),
  FUNCTION 7 g_int_same (_int4, _int4, internal);
```

Compatibility

CREATE OPERATOR CLASS is a PostgreSQL extension. There is no CREATE OPERATOR CLASS statement in the SQL standard.

See Also

ALTER OPERATOR CLASS, DROP OPERATOR CLASS, CREATE OPERATOR FAMILY, ALTER OPERATOR FAMILY
CREATE OPERATOR FAMILY

CREATE OPERATOR FAMILY — define a new operator family

Synopsis

CREATE OPERATOR FAMILY name USING index_method

Description

CREATE OPERATOR FAMILY creates a new operator family. An operator family defines a collection of related operator classes, and perhaps some additional operators and support functions that are compatible with these operator classes but not essential for the functioning of any individual index. (Operators and functions that are essential to indexes should be grouped within the relevant operator class, rather than being "loose" in the operator family. Typically, single-data-type operators are bound to operator classes, while cross-data-type operators can be loose in an operator family containing operator classes for both data types.)

The new operator family is initially empty. It should be populated by issuing subsequent CREATE OPERATOR CLASS commands to add contained operator classes, and optionally ALTER OPERATOR FAMILY commands to add “loose” operators and their corresponding support functions.

If a schema name is given then the operator family is created in the specified schema. Otherwise it is created in the current schema. Two operator families in the same schema can have the same name only if they are for different index methods.

The user who defines an operator family becomes its owner. Presently, the creating user must be a superuser. (This restriction is made because an erroneous operator family definition could confuse or even crash the server.)

Refer to Section 38.15 for further information.

Parameters

name

The name of the operator family to be created. The name can be schema-qualified.

index_method

The name of the index method this operator family is for.

Compatibility

CREATE OPERATOR FAMILY is a PostgreSQL extension. There is no CREATE OPERATOR FAMILY statement in the SQL standard.

See Also

ALTER OPERATOR FAMILY, DROP OPERATOR FAMILY, CREATE OPERATOR CLASS, ALTER OPERATOR CLASS, DROP OPERATOR CLASS
CREATE POLICY

CREATE POLICY — define a new row level security policy for a table

Synopsis

CREATE POLICY name ON table_name
[ AS { PERMISSIVE | RESTRICTIVE } ]
[ FOR { ALL | SELECT | INSERT | UPDATE | DELETE } ]
[ TO { role_name | PUBLIC | CURRENT_USER | SESSION_USER } [, ... ] ]
[ USING ( using_expression ) ]
[ WITH CHECK ( check_expression ) ]

Description

The CREATE POLICY command defines a new row-level security policy for a table. Note that row-level security must be enabled on the table (using ALTER TABLE ... ENABLE ROW LEVEL SECURITY) in order for created policies to be applied.

A policy grants the permission to select, insert, update, or delete rows that match the relevant policy expression. Existing table rows are checked against the expression specified in USING, while new rows that would be created via INSERT or UPDATE are checked against the expression specified in WITH CHECK. When a USING expression returns true for a given row then that row is visible to the user, while if false or null is returned then the row is not visible. When a WITH CHECK expression returns true for a row then that row is inserted or updated, while if false or null is returned then an error occurs.

For INSERT and UPDATE statements, WITH CHECK expressions are enforced after BEFORE triggers are fired, and before any actual data modifications are made. Thus a BEFORE ROW trigger may modify the data to be inserted, affecting the result of the security policy check. WITH CHECK expressions are enforced before any other constraints.

Policy names are per-table. Therefore, one policy name can be used for many different tables and have a definition for each table which is appropriate to that table.

Policies can be applied for specific commands or for specific roles. The default for newly created policies is that they apply for all commands and roles, unless otherwise specified. Multiple policies may apply to a single command; see below for more details. Table 240 summarizes how the different types of policy apply to specific commands.

For policies that can have both USING and WITH CHECK expressions (ALL and UPDATE), if no WITH CHECK expression is defined, then the USING expression will be used both to determine which rows are visible (normal USING case) and which new rows will be allowed to be added (WITH CHECK case).

If row-level security is enabled for a table, but no applicable policies exist, a “default deny” policy is assumed, so that no rows will be visible or updatable.

Parameters

\( name \)

The name of the policy to be created. This must be distinct from the name of any other policy for the table.
CREATE POLICY

table_name

The name (optionally schema-qualified) of the table the policy applies to.

PERMISSIVE

Specify that the policy is to be created as a permissive policy. All permissive policies which are applicable to a given query will be combined together using the Boolean “OR” operator. By creating permissive policies, administrators can add to the set of records which can be accessed. Policies are permissive by default.

RESTRICTIVE

Specify that the policy is to be created as a restrictive policy. All restrictive policies which are applicable to a given query will be combined together using the Boolean “AND” operator. By creating restrictive policies, administrators can reduce the set of records which can be accessed as all restrictive policies must be passed for each record.

Note that there needs to be at least one permissive policy to grant access to records before restrictive policies can be usefully used to reduce that access. If only restrictive policies exist, then no records will be accessible. When a mix of permissive and restrictive policies are present, a record is only accessible if at least one of the permissive policies passes, in addition to all the restrictive policies.

command

The command to which the policy applies. Valid options are ALL, SELECT, INSERT, UPDATE, and DELETE. ALL is the default. See below for specifics regarding how these are applied.

role_name

The role(s) to which the policy is to be applied. The default is PUBLIC, which will apply the policy to all roles.

using_expression

Any SQL conditional expression (returning boolean). The conditional expression cannot contain any aggregate or window functions. This expression will be added to queries that refer to the table if row level security is enabled. Rows for which the expression returns true will be visible. Any rows for which the expression returns false or null will not be visible to the user (in a SELECT), and will not be available for modification (in an UPDATE or DELETE). Such rows are silently suppressed; no error is reported.

check_expression

Any SQL conditional expression (returning boolean). The conditional expression cannot contain any aggregate or window functions. This expression will be used in INSERT and UPDATE queries against the table if row level security is enabled. Only rows for which the expression evaluates to true will be allowed. An error will be thrown if the expression evaluates to false or null for any of the records inserted or any of the records that result from the update. Note that the check_expression is evaluated against the proposed new contents of the row, not the original contents.

Per-Command Policies

ALL

Using ALL for a policy means that it will apply to all commands, regardless of the type of command. If an ALL policy exists and more specific policies exist, then both the ALL policy and the more specific policy (or policies) will be applied. Additionally, ALL policies will be applied to
CREATE POLICY

both the selection side of a query and the modification side, using the USING expression for both cases if only a USING expression has been defined.

As an example, if an UPDATE is issued, then the ALL policy will be applicable both to what the UPDATE will be able to select as rows to be updated (applying the USING expression), and to the resulting updated rows, to check if they are permitted to be added to the table (applying the WITH CHECK expression, if defined, and the USING expression otherwise). If an INSERT or UPDATE command attempts to add rows to the table that do not pass the ALL policy’s WITH CHECK expression, the entire command will be aborted.

SELECT

Using SELECT for a policy means that it will apply to SELECT queries and whenever SELECT permissions are required on the relation the policy is defined for. The result is that only those records from the relation that pass the SELECT policy will be returned during a SELECT query, and that queries that require SELECT permissions, such as UPDATE, will also only see those records that are allowed by the SELECT policy. A SELECT policy cannot have a WITH CHECK expression, as it only applies in cases where records are being retrieved from the relation.

INSERT

Using INSERT for a policy means that it will apply to INSERT commands. Rows being inserted that do not pass this policy will result in a policy violation error, and the entire INSERT command will be aborted. An INSERT policy cannot have a USING expression, as it only applies in cases where records are being added to the relation.

Note that INSERT with ON CONFLICT DO UPDATE checks INSERT policies’ WITH CHECK expressions only for rows appended to the relation by the INSERT path.

UPDATE

Using UPDATE for a policy means that it will apply to UPDATE, SELECT FOR UPDATE and SELECT FOR SHARE commands, as well as auxiliary ON CONFLICT DO UPDATE clauses of INSERT commands. Since UPDATE involves pulling an existing record and replacing it with a new modified record, UPDATE policies accept both a USING expression and a WITH CHECK expression. The USING expression determines which records the UPDATE command will see to operate against, while the WITH CHECK expression defines which modified rows are allowed to be stored back into the relation.

Any rows whose updated values do not pass the WITH CHECK expression will cause an error, and the entire command will be aborted. If only a USING clause is specified, then that clause will be used for both USING and WITH CHECK cases.

Typically an UPDATE command also needs to read data from columns in the relation being updated (e.g., in a WHERE clause or a RETURNING clause, or in an expression on the right hand side of the SET clause). In this case, SELECT rights are also required on the relation being updated, and the appropriate SELECT or ALL policies will be applied in addition to the UPDATE policies. Thus the user must have access to the row(s) being updated through a SELECT or ALL policy in addition to being granted permission to update the row(s) via an UPDATE or ALL policy.

When an INSERT command has an auxiliary ON CONFLICT DO UPDATE clause, if the UPDATE path is taken, the row to be updated is first checked against the USING expressions of any UPDATE policies, and then the new updated row is checked against the WITH CHECK expressions. Note, however, that unlike a standalone UPDATE command, if the existing row does not pass the USING expressions, an error will be thrown (the UPDATE path will never be silently avoided).

DELETE

Using DELETE for a policy means that it will apply to DELETE commands. Only rows that pass this policy will be seen by a DELETE command. There can be rows that are visible through a
SELECT that are not available for deletion, if they do not pass the USING expression for the DELETE policy.

In most cases a DELETE command also needs to read data from columns in the relation that it is deleting from (e.g., in a WHERE clause or a RETURNING clause). In this case, SELECT rights are also required on the relation, and the appropriate SELECT or ALL policies will be applied in addition to the DELETE policies. Thus the user must have access to the row(s) being deleted through a SELECT or ALL policy in addition to being granted permission to delete the row(s) via a DELETE or ALL policy.

A DELETE policy cannot have a WITH CHECK expression, as it only applies in cases where records are being deleted from the relation, so that there is no new row to check.

### Table 240. Policies Applied by Command Type

<table>
<thead>
<tr>
<th>Command</th>
<th>SELECT/ALL policy</th>
<th>INSERT/ALL policy</th>
<th>UPDATE/ALL policy</th>
<th>DELETE/ALL policy</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>USING expression</td>
<td>WITH CHECK expression</td>
<td>USING expression</td>
<td>WITH CHECK expression</td>
</tr>
<tr>
<td>SELECT</td>
<td>Existing row</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>SELECT FOR UPDATE/SHARE</td>
<td>Existing row</td>
<td>—</td>
<td>Existing row</td>
<td>—</td>
</tr>
<tr>
<td>INSERT</td>
<td>—</td>
<td>New row</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>INSERT ... RETURNING</td>
<td>New row</td>
<td>New row</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>UPDATE</td>
<td>Existing &amp; new rows</td>
<td>—</td>
<td>Existing row</td>
<td>New row</td>
</tr>
<tr>
<td>DELETE</td>
<td>Existing row</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>ON CONFLICT DO UPDATE</td>
<td>Existing &amp; new rows</td>
<td>—</td>
<td>Existing row</td>
<td>New row</td>
</tr>
</tbody>
</table>

* If read access is required to the existing or new row (for example, a WHERE or RETURNING clause that refers to columns from the relation).

### Application of Multiple Policies

When multiple policies of different command types apply to the same command (for example, SELECT and UPDATE policies applied to an UPDATE command), then the user must have both types of permissions (for example, permission to select rows from the relation as well as permission to update them). Thus the expressions for one type of policy are combined with the expressions for the other type of policy using the AND operator.

When multiple policies of the same command type apply to the same command, then there must be at least one PERMISSIVE policy granting access to the relation, and all of the RESTRICTIVE policies must pass. Thus all the PERMISSIVE policy expressions are combined using OR, all the RESTRICTIVE policy expressions are combined using AND, and the results are combined using AND. If there are no PERMISSIVE policies, then access is denied.

Note that, for the purposes of combining multiple policies, ALL policies are treated as having the same type as whichever other type of policy is being applied.

For example, in an UPDATE command requiring both SELECT and UPDATE permissions, if there are multiple applicable policies of each type, they will be combined as follows:
**CREATE POLICY**

(expression from RESTRICTIVE SELECT/ALL policy 1
 AND
 expression from RESTRICTIVE SELECT/ALL policy 2
 AND
 ... 
 AND
 (expression from PERMISSIVE SELECT/ALL policy 1
 OR
 expression from PERMISSIVE SELECT/ALL policy 2
 OR
 ... 
)
 AND
 expression from RESTRICTIVE UPDATE/ALL policy 1
 AND
 expression from RESTRICTIVE UPDATE/ALL policy 2
 AND
 ... 
 AND
 (expression from PERMISSIVE UPDATE/ALL policy 1
 OR
 expression from PERMISSIVE UPDATE/ALL policy 2
 OR
 ... 
)

**Notes**

You must be the owner of a table to create or change policies for it.

While policies will be applied for explicit queries against tables in the database, they are not applied when the system is performing internal referential integrity checks or validating constraints. This means there are indirect ways to determine that a given value exists. An example of this is attempting to insert a duplicate value into a column that is a primary key or has a unique constraint. If the insert fails then the user can infer that the value already exists. (This example assumes that the user is permitted by policy to insert records which they are not allowed to see.) Another example is where a user is allowed to insert into a table which references another, otherwise hidden table. Existence can be determined by the user inserting values into the referencing table, where success would indicate that the value exists in the referenced table. These issues can be addressed by carefully crafting policies to prevent users from being able to insert, delete, or update records at all which might possibly indicate a value they are not otherwise able to see, or by using generated values (e.g., surrogate keys) instead of keys with external meanings.

Generally, the system will enforce filter conditions imposed using security policies prior to qualifications that appear in user queries, in order to prevent inadvertent exposure of the protected data to user-defined functions which might not be trustworthy. However, functions and operators marked by the system (or the system administrator) as LEAKPROOF may be evaluated before policy expressions, as they are assumed to be trustworthy.

Since policy expressions are added to the user’s query directly, they will be run with the rights of the user running the overall query. Therefore, users who are using a given policy must be able to access any tables or functions referenced in the expression or they will simply receive a permission denied error when attempting to query the table that has row-level security enabled. This does not change how views work, however. As with normal queries and views, permission checks and policies for the tables which are referenced by a view will use the view owner’s rights and any policies which apply to the view owner.
Additional discussion and practical examples can be found in Section 5.7.

Compatibility

CREATE POLICY is a PostgreSQL extension.

See Also

ALTER POLICY, DROP POLICY, ALTER TABLE
CREATE PROCEDURE

CREATE PROCEDURE — define a new procedure

Synopsis

CREATE [ OR REPLACE ] PROCEDURE
   name ( [ [ argmode ] [ argname ] argtype [ { DEFAULT | = } default_expr ] [, ...] ] )
   { LANGUAGE lang_name
     | TRANSFORM { FOR TYPE type_name } [, ... ]
     | [ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
     | SET configuration_parameter { TO value | = value | FROM CURRENT }
     | AS 'definition'
     | AS 'obj_file', 'link_symbol'
   ) ...

Description

CREATE PROCEDURE defines a new procedure. CREATE OR REPLACE PROCEDURE will either create a new procedure, or replace an existing definition. To be able to define a procedure, the user must have the USAGE privilege on the language.

If a schema name is included, then the procedure is created in the specified schema. Otherwise it is created in the current schema. The name of the new procedure must not match any existing procedure or function with the same input argument types in the same schema. However, procedures and functions of different argument types can share a name (this is called overloading).

To replace the current definition of an existing procedure, use CREATE OR REPLACE PROCEDURE. It is not possible to change the name or argument types of a procedure this way (if you tried, you would actually be creating a new, distinct procedure).

When CREATE OR REPLACE PROCEDURE is used to replace an existing procedure, the ownership and permissions of the procedure do not change. All other procedure properties are assigned the values specified or implied in the command. You must own the procedure to replace it (this includes being a member of the owning role).

The user that creates the procedure becomes the owner of the procedure.

To be able to create a procedure, you must have USAGE privilege on the argument types.

Parameters

name

The name (optionally schema-qualified) of the procedure to create.

argmode

The mode of an argument: IN, INOUT, or VARIADIC. If omitted, the default is IN. (OUT arguments are currently not supported for procedures. Use INOUT instead.)

argname

The name of an argument.
CREATE PROCEDURE

argtype

The data type(s) of the procedure's arguments (optionally schema-qualified), if any. The argument
types can be base, composite, or domain types, or can reference the type of a table column.

Depending on the implementation language it might also be allowed to specify “pseudo-types”
such as cstring. Pseudo-types indicate that the actual argument type is either incompletely
specified, or outside the set of ordinary SQL data types.

The type of a column is referenced by writing table_name.column_name%TYPE. Using this
feature can sometimes help make a procedure independent of changes to the definition of a table.

default_expr

An expression to be used as default value if the parameter is not specified. The expression has to
be coercible to the argument type of the parameter. All input parameters following a parameter
with a default value must have default values as well.

lang_name

The name of the language that the procedure is implemented in. It can be sql, c, internal,
or the name of a user-defined procedural language, e.g. plpgsql. Enclosing the name in single
quotes is deprecated and requires matching case.

TRANSFORM { FOR TYPE type_name } [, ... ] }

Lists which transforms a call to the procedure should apply. Transforms convert between SQL
types and language-specific data types; see CREATE TRANSFORM. Procedural language im-
plementations usually have hardcoded knowledge of the built-in types, so those don't need to be
listed here. If a procedural language implementation does not know how to handle a type and no
transform is supplied, it will fall back to a default behavior for converting data types, but this
depends on the implementation.

[EXTERNAL] SECURITY INVOKER
[EXTERNAL] SECURITY DEFINER

SECURITY INVOKER indicates that the procedure is to be executed with the privileges of the
user that calls it. That is the default. SECURITY DEFINER specifies that the procedure is to be
executed with the privileges of the user that owns it.

The key word EXTERNAL is allowed for SQL conformance, but it is optional since, unlike in
SQL, this feature applies to all procedures not only external ones.

A SECURITY DEFINER procedure cannot execute transaction control statements (for example,
COMMIT and ROLLBACK, depending on the language).

configuration_parameter
value

The SET clause causes the specified configuration parameter to be set to the specified value
when the procedure is entered, and then restored to its prior value when the procedure exits. SET
FROM CURRENT saves the value of the parameter that is current when CREATE PROCEDURE
is executed as the value to be applied when the procedure is entered.

If a SET clause is attached to a procedure, then the effects of a SET LOCAL command execut-
ed inside the procedure for the same variable are restricted to the procedure: the configuration
parameter's prior value is still restored at procedure exit. However, an ordinary SET command
(without LOCAL) overrides the SET clause, much as it would do for a previous SET LOCAL
command: the effects of such a command will persist after procedure exit, unless the current
transaction is rolled back.
CREATE PROCEDURE

If a SET clause is attached to a procedure, then that procedure cannot execute transaction control statements (for example, COMMIT and ROLLBACK, depending on the language).

See SET and Chapter 19 for more information about allowed parameter names and values.

definition

A string constant defining the procedure; the meaning depends on the language. It can be an internal procedure name, the path to an object file, an SQL command, or text in a procedural language.

It is often helpful to use dollar quoting (see Section 4.1.2.4) to write the procedure definition string, rather than the normal single quote syntax. Without dollar quoting, any single quotes or backslashes in the procedure definition must be escaped by doubling them.

obj_file, link_symbol

This form of the AS clause is used for dynamically loadable C language procedures when the procedure name in the C language source code is not the same as the name of the SQL procedure. The string obj_file is the name of the shared library file containing the compiled C procedure, and is interpreted as for the LOAD command. The string link_symbol is the procedure's link symbol, that is, the name of the procedure in the C language source code. If the link symbol is omitted, it is assumed to be the same as the name of the SQL procedure being defined.

When repeated CREATE PROCEDURE calls refer to the same object file, the file is only loaded once per session. To unload and reload the file (perhaps during development), start a new session.

Notes

See CREATE FUNCTION for more details on function creation that also apply to procedures.

Use CALL to execute a procedure.

Examples

CREATE PROCEDURE insert_data(a integer, b integer)
LANGUAGE SQL
AS $$
INSERT INTO tbl VALUES (a);
INSERT INTO tbl VALUES (b);
$$;

CALL insert_data(1, 2);

Compatibility

A CREATE PROCEDURE command is defined in the SQL standard. The PostgreSQL version is similar but not fully compatible. For details see also CREATE FUNCTION.

See Also

ALTER PROCEDURE, DROP PROCEDURE, CALL, CREATE FUNCTION
CREATE PUBLICATION

CREATE PUBLICATION — define a new publication

Synopsis

CREATE PUBLICATION name
   [ FOR TABLE [ ONLY ] table_name [ * ] [, ... ]
   | FOR ALL TABLES ]
   [ WITH ( publication_parameter [= value] [, ... ] ) ]

Description

CREATE PUBLICATION adds a new publication into the current database. The publication name must be distinct from the name of any existing publication in the current database.

A publication is essentially a group of tables whose data changes are intended to be replicated through logical replication. See Section 31.1 for details about how publications fit into the logical replication setup.

Parameters

name

The name of the new publication.

FOR TABLE

Specifies a list of tables to add to the publication. If ONLY is specified before the table name, only that table is added to the publication. If ONLY is not specified, the table and all its descendant tables (if any) are added. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

Only persistent base tables can be part of a publication. Temporary tables, unlogged tables, foreign tables, materialized views, regular views, and partitioned tables cannot be part of a publication. To replicate a partitioned table, add the individual partitions to the publication.

FOR ALL TABLES

Marks the publication as one that replicates changes for all tables in the database, including tables created in the future.

WITH ( publication_parameter [= value] [, ... ] )

This clause specifies optional parameters for a publication. The following parameters are supported:

publish(string)

This parameter determines which DML operations will be published by the new publication to the subscribers. The value is comma-separated list of operations. The allowed operations are insert, update, delete, and truncate. The default is to publish all actions, and so the default value for this option is 'insert, update, delete, truncate'.
CREATE PUBLICATION

Notes

If neither FOR TABLE nor FOR ALL TABLES is specified, then the publication starts out with an empty set of tables. That is useful if tables are to be added later.

The creation of a publication does not start replication. It only defines a grouping and filtering logic for future subscribers.

To create a publication, the invoking user must have the CREATE privilege for the current database. (Of course, superusers bypass this check.)

To add a table to a publication, the invoking user must have ownership rights on the table. The FOR ALL TABLES clause requires the invoking user to be a superuser.

The tables added to a publication that publishes UPDATE and/or DELETE operations must have REPLICA IDENTITY defined. Otherwise those operations will be disallowed on those tables.

For an INSERT ... ON CONFLICT command, the publication will publish the operation that actually results from the command. So depending of the outcome, it may be published as either INSERT or UPDATE, or it may not be published at all.

COPY ... FROM commands are published as INSERT operations.

DDL operations are not published.

Examples

Create a publication that publishes all changes in two tables:

CREATE PUBLICATION mypublication FOR TABLE users, departments;

Create a publication that publishes all changes in all tables:

CREATE PUBLICATION alltables FOR ALL TABLES;

Create a publication that only publishes INSERT operations in one table:

CREATE PUBLICATION insert_only FOR TABLE mydata
   WITH (publish = 'insert');

Compatibility

CREATE PUBLICATION is a PostgreSQL extension.

See Also

ALTER PUBLICATION, DROP PUBLICATION
CREATE ROLE

CREATE ROLE — define a new database role

Synopsis

CREATE ROLE name [ [ WITH ] option [ ... ] ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICATION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT connlimit
- [ ENCRYPTED ] PASSWORD 'password'
- VALID UNTIL 'timestamp'
- IN ROLE role_name [, ...]
- IN GROUP role_name [, ...]
- ROLE role_name [, ...]
- ADMIN role_name [, ...]
- USER role_name [, ...]
- SYSID uid

Description

CREATE ROLE adds a new role to a PostgreSQL database cluster. A role is an entity that can own database objects and have database privileges; a role can be considered a “user”, a “group”, or both depending on how it is used. Refer to Chapter 21 and Chapter 20 for information about managing users and authentication. You must have CREATEROLE privilege or be a database superuser to use this command.

Note that roles are defined at the database cluster level, and so are valid in all databases in the cluster.

Parameters

name

The name of the new role.

SUPERUSER
NOSUPERUSER

These clauses determine whether the new role is a “superuser”, who can override all access restrictions within the database. Superuser status is dangerous and should be used only when really needed. You must yourself be a superuser to create a new superuser. If not specified, NOSUPERUSER is the default.
CREATE ROLE

CREATEDB
NOCREATEDB

These clauses define a role's ability to create databases. If CREATEDB is specified, the role being defined will be allowed to create new databases. Specifying NOCREATEDB will deny a role the ability to create databases. If not specified, NOCREATEDB is the default.

CREATEROLE
NOCREATEROLE

These clauses determine whether a role will be permitted to create new roles (that is, execute CREATE ROLE). A role with CREATEROLE privilege can also alter and drop other roles. If not specified, NOCREATEROLE is the default.

INHERIT
NOINHERIT

These clauses determine whether a role “inherits” the privileges of roles it is a member of. A role with the INHERIT attribute can automatically use whatever database privileges have been granted to all roles it is directly or indirectly a member of. Without INHERIT, membership in another role only grants the ability to SET ROLE to that other role; the privileges of the other role are only available after having done so. If not specified, INHERIT is the default.

LOGIN
NOLOGIN

These clauses determine whether a role is allowed to log in; that is, whether the role can be given as the initial session authorization name during client connection. A role having the LOGIN attribute can be thought of as a user. Roles without this attribute are useful for managing database privileges, but are not users in the usual sense of the word. If not specified, NOLOGIN is the default, except when CREATE ROLE is invoked through its alternative spelling CREATE USER.

REPLICATION
NOREPLICATION

These clauses determine whether a role is a replication role. A role must have this attribute (or be a superuser) in order to be able to connect to the server in replication mode (physical or logical replication) and in order to be able to create or drop replication slots. A role having the REPLICATION attribute is a very highly privileged role, and should only be used on roles actually used for replication. If not specified, NOREPLICATION is the default.

BYPASSRLS
NOBYPASSRLS

These clauses determine whether a role bypasses every row-level security (RLS) policy. NOBYPASSRLS is the default. Note that pg_dump will set row_security to OFF by default, to ensure all contents of a table are dumped out. If the user running pg_dump does not have appropriate permissions, an error will be returned. The superuser and owner of the table being dumped always bypass RLS.

CONNECTION LIMIT connlimit

If role can log in, this specifies how many concurrent connections the role can make. -1 (the default) means no limit. Note that only normal connections are counted towards this limit. Neither prepared transactions nor background worker connections are counted towards this limit.

[ ENCRYPTED ] PASSWORD password

Sets the role's password. (A password is only of use for roles having the LOGIN attribute, but you can nonetheless define one for roles without it.) If you do not plan to use password authentication
you can omit this option. If no password is specified, the password will be set to null and password authentication will always fail for that user. A null password can optionally be written explicitly as PASSWORD NULL.

**Note**

Specifying an empty string will also set the password to null, but that was not the case before PostgreSQL version 10. In earlier versions, an empty string could be used, or not, depending on the authentication method and the exact version, and libpq would refuse to use it in any case. To avoid the ambiguity, specifying an empty string should be avoided.

The password is always stored encrypted in the system catalogs. The ENCRYPTED keyword has no effect, but is accepted for backwards compatibility. The method of encryption is determined by the configuration parameter password_encryption. If the presented password string is already in MD5-encrypted or SCRAM-encrypted format, then it is stored as-is regardless of password_encryption (since the system cannot decrypt the specified encrypted password string, to encrypt it in a different format). This allows reloading of encrypted passwords during dump/restore.

**VALID UNTIL 'timestamp'**

The VALID UNTIL clause sets a date and time after which the role's password is no longer valid. If this clause is omitted the password will be valid for all time.

**IN ROLE role_name**

The IN ROLE clause lists one or more existing roles to which the new role will be immediately added as a new member. (Note that there is no option to add the new role as an administrator; use a separate GRANT command to do that.)

**IN GROUP role_name**

IN GROUP is an obsolete spelling of IN ROLE.

**ROLE role_name**

The ROLE clause lists one or more existing roles which are automatically added as members of the new role. (This in effect makes the new role a "group".)

**ADMIN role_name**

The ADMIN clause is like ROLE, but the named roles are added to the new role WITH ADMIN OPTION, giving them the right to grant membership in this role to others.

**USER role_name**

The USER clause is an obsolete spelling of the ROLE clause.

**SYSID uid**

The SYSID clause is ignored, but is accepted for backwards compatibility.

**Notes**

Use ALTER ROLE to change the attributes of a role, and DROP ROLE to remove a role. All the attributes specified by CREATE ROLE can be modified by later ALTER ROLE commands.
The preferred way to add and remove members of roles that are being used as groups is to use Grant
and Revoke.

The **VALID UNTIL** clause defines an expiration time for a password only, not for the role per se. In
particular, the expiration time is not enforced when logging in using a non-password-based authenti-
cation method.

The **INHERIT** attribute governs inheritance of grantable privileges (that is, access privileges for data-
base objects and role memberships). It does not apply to the special role attributes set by CREATE
ROLE and ALTER ROLE. For example, being a member of a role with **CREATEDB** privilege does not
immediately grant the ability to create databases, even if **INHERIT** is set; it would be necessary to
become that role via **SET ROLE** before creating a database.

The **INHERIT** attribute is the default for reasons of backwards compatibility: in prior releases of
PostgreSQL, users always had access to all privileges of groups they were members of. However,
**NOINHERIT** provides a closer match to the semantics specified in the SQL standard.

Be careful with the **CREATEROLE** privilege. There is no concept of inheritance for the privileges of
a **CREATEROLE**-role. That means that even if a role does not have a certain privilege but is allowed
to create other roles, it can easily create another role with different privileges than its own (except
for creating roles with superuser privileges). For example, if the role “user” has the **CREATEROLE**
privilege but not the **CREATEDB** privilege, nonetheless it can create a new role with the **CREATEDB**
privilege. Therefore, regard roles that have the **CREATEROLE** privilege as almost-superuser-roles.

PostgreSQL includes a program **createuser** that has the same functionality as **CREATE ROLE** (in fact,
it calls this command) but can be run from the command shell.

The **CONNECTION LIMIT** option is only enforced approximately; if two new sessions start at about
the same time when just one connection “slot” remains for the role, it is possible that both will fail.
Also, the limit is never enforced for superusers.

Caution must be exercised when specifying an unencrypted password with this command. The pass-
word will be transmitted to the server in cleartext, and it might also be logged in the client’s command
history or the server log. The command **createuser**, however, transmits the password encrypted. Also,
**psql** contains a command **\password** that can be used to safely change the password later.

**Examples**

Create a role that can log in, but don't give it a password:

```
CREATE ROLE jonathan LOGIN;
```

Create a role with a password:

```
CREATE USER davide WITH PASSWORD 'jw8s0F4';
```

(CREATE USER is the same as CREATE ROLE except that it implies LOGIN.)

Create a role with a password that is valid until the end of 2004. After one second has ticked in 2005,
the password is no longer valid.

```
CREATE ROLE miriam WITH LOGIN PASSWORD 'jw8s0F4' VALID UNTIL '2005-01-01';
```

Create a role that can create databases and manage roles:
CREATE ROLE admin WITH CREATEDB CREATEROLE;

Compatibility

The CREATE ROLE statement is in the SQL standard, but the standard only requires the syntax

CREATE ROLE name [ WITH ADMIN role_name ]

Multiple initial administrators, and all the other options of CREATE ROLE, are PostgreSQL extensions.

The SQL standard defines the concepts of users and roles, but it regards them as distinct concepts and leaves all commands defining users to be specified by each database implementation. In PostgreSQL we have chosen to unify users and roles into a single kind of entity. Roles therefore have many more optional attributes than they do in the standard.

The behavior specified by the SQL standard is most closely approximated by giving users the NOINHERIT attribute, while roles are given the INHERIT attribute.

See Also

SET ROLE, ALTER ROLE, DROP ROLE, GRANT, REVOKE, createuser
CREATE RULE

CREATE RULE — define a new rewrite rule

Synopsis

CREATE [ OR REPLACE ] RULE name AS ON event
TO table_name [ WHERE condition ]
DO [ ALSO | INSTEAD ] { NOTHING | command | ( command ; command ...

where event can be one of:

SELECT | INSERT | UPDATE | DELETE

Description

CREATE RULE defines a new rule applying to a specified table or view. CREATE OR REPLACE RULE will either create a new rule, or replace an existing rule of the same name for the same table.

The PostgreSQL rule system allows one to define an alternative action to be performed on insertions, updates, or deletions in database tables. Roughly speaking, a rule causes additional commands to be executed when a given command on a given table is executed. Alternatively, an INSTEAD rule can replace a given command by another, or cause a command not to be executed at all. Rules are used to implement SQL views as well. It is important to realize that a rule is really a command transformation mechanism, or command macro. The transformation happens before the execution of the command starts. If you actually want an operation that fires independently for each physical row, you probably want to use a trigger, not a rule. More information about the rules system is in Chapter 41.

Presently, ON SELECT rules must be unconditional INSTEAD rules and must have actions that consist of a single SELECT command. Thus, an ON SELECT rule effectively turns the table into a view, whose visible contents are the rows returned by the rule's SELECT command rather than whatever had been stored in the table (if anything). It is considered better style to write a CREATE VIEW command than to create a real table and define an ON SELECT rule for it.

You can create the illusion of an updatable view by defining ON INSERT, ON UPDATE, and ON DELETE rules (or any subset of those that's sufficient for your purposes) to replace update actions on the view with appropriate updates on other tables. If you want to support INSERT RETURNING and so on, then be sure to put a suitable RETURNING clause into each of these rules.

There is a catch if you try to use conditional rules for complex view updates: there must be an unconditional INSTEAD rule for each action you wish to allow on the view. If the rule is conditional, or is not INSTEAD, then the system will still reject attempts to perform the update action, because it thinks it might end up trying to perform the action on the dummy table of the view in some cases. If you want to handle all the useful cases in conditional rules, add an unconditional DO INSTEAD NOTHING rule to ensure that the system understands it will never be called on to update the dummy table. Then make the conditional rules non-INSTEAD; in the cases where they are applied, they add to the default INSTEAD NOTHING action. (This method does not currently work to support RETURNING queries, however.)

Note

A view that is simple enough to be automatically updatable (see CREATE VIEW) does not require a user-created rule in order to be updatable. While you can create an
explicit rule anyway, the automatic update transformation will generally outperform an explicit rule.

Another alternative worth considering is to use INSTEAD OF triggers (see CREATE TRIGGER) in place of rules.

### Parameters

**name**

The name of a rule to create. This must be distinct from the name of any other rule for the same table. Multiple rules on the same table and same event type are applied in alphabetical name order.

**event**

The event is one of SELECT, INSERT, UPDATE, or DELETE. Note that an INSERT containing an ON CONFLICT clause cannot be used on tables that have either INSERT or UPDATE rules. Consider using an updatable view instead.

**table_name**

The name (optionally schema-qualified) of the table or view the rule applies to.

**condition**

Any SQL conditional expression (returning boolean). The condition expression cannot refer to any tables except NEW and OLD, and cannot contain aggregate functions.

**INSTEAD**

**INSTEAD** indicates that the commands should be executed instead of the original command.

**ALSO**

**ALSO** indicates that the commands should be executed in addition to the original command.

If neither ALSO nor INSTEAD is specified, ALSO is the default.

**command**

The command or commands that make up the rule action. Valid commands are SELECT, INSERT, UPDATE, DELETE, or NOTIFY.

Within condition and command, the special table names NEW and OLD can be used to refer to values in the referenced table. NEW is valid in ON INSERT and ON UPDATE rules to refer to the new row being inserted or updated. OLD is valid in ON UPDATE and ON DELETE rules to refer to the existing row being updated or deleted.

### Notes

You must be the owner of a table to create or change rules for it.

In a rule for INSERT, UPDATE, or DELETE on a view, you can add a RETURNING clause that emits the view's columns. This clause will be used to compute the outputs if the rule is triggered by an INSERT RETURNING, UPDATE RETURNING, or DELETE RETURNING command respectively. When the rule is triggered by a command without RETURNING, the rule's RETURNING clause will be ignored. The current implementation allows only unconditional INSTEAD rules to contain RETURNING; furthermore there can be at most one RETURNING clause among all the rules for the same event. (This ensures that there is only one candidate RETURNING clause to be used to compute
the results.) RETURNING queries on the view will be rejected if there is no RETURNING clause in any available rule.

It is very important to take care to avoid circular rules. For example, though each of the following two rule definitions are accepted by PostgreSQL, the SELECT command would cause PostgreSQL to report an error because of recursive expansion of a rule:

```sql
CREATE RULE "_RETURN" AS 
  ON SELECT TO t1 
  DO INSTEAD 
    SELECT * FROM t2;

CREATE RULE "_RETURN" AS 
  ON SELECT TO t2 
  DO INSTEAD 
    SELECT * FROM t1;

SELECT * FROM t1;
```

Presently, if a rule action contains a NOTIFY command, the NOTIFY command will be executed unconditionally, that is, the NOTIFY will be issued even if there are not any rows that the rule should apply to. For example, in:

```sql
CREATE RULE notify_me AS ON UPDATE TO mytable DO ALSO NOTIFY 
  mytable;

UPDATE mytable SET name = 'foo' WHERE id = 42;
```

one NOTIFY event will be sent during the UPDATE, whether or not there are any rows that match the condition $id = 42$. This is an implementation restriction that might be fixed in future releases.

**Compatibility**

CREATE RULE is a PostgreSQL language extension, as is the entire query rewrite system.

**See Also**

ALERT RULE, DROP RULE
CREATE SCHEMA

CREATE SCHEMA — define a new schema

Synopsis

CREATE SCHEMA schema_name [ AUTHORIZATION role_specification ] [ schema_element [ ... ] ]
CREATE SCHEMA AUTHORIZATION role_specification [ schema_element [ ... ] ]
CREATE SCHEMA IF NOT EXISTS schema_name
CREATE SCHEMA IF NOT EXISTS AUTHORIZATION role_specification

where role_specification can be:

user_name
| CURRENT_USER
| SESSION_USER

Description

CREATE SCHEMA enters a new schema into the current database. The schema name must be distinct from the name of any existing schema in the current database.

A schema is essentially a namespace: it contains named objects (tables, data types, functions, and operators) whose names can duplicate those of other objects existing in other schemas. Named objects are accessed either by “qualifying” their names with the schema name as a prefix, or by setting a search path that includes the desired schema(s). A CREATE command specifying an unqualified object name creates the object in the current schema (the one at the front of the search path, which can be determined with the function current_schema).

Optionally, CREATE SCHEMA can include subcommands to create objects within the new schema. The subcommands are treated essentially the same as separate commands issued after creating the schema, except that if the AUTHORIZATION clause is used, all the created objects will be owned by that user.

Parameters

schema_name

The name of a schema to be created. If this is omitted, the user_name is used as the schema name. The name cannot begin with pg_, as such names are reserved for system schemas.

user_name

The role name of the user who will own the new schema. If omitted, defaults to the user executing the command. To create a schema owned by another role, you must be a direct or indirect member of that role, or be a superuser.

schema_element

An SQL statement defining an object to be created within the schema. Currently, only CREATE TABLE, CREATE VIEW, CREATE INDEX, CREATE SEQUENCE, CREATE TRIGGER and
CREATE SCHEMA

GRANT are accepted as clauses within CREATE SCHEMA. Other kinds of objects may be created in separate commands after the schema is created.

IF NOT EXISTS

Do nothing (except issuing a notice) if a schema with the same name already exists. schema_element subcommands cannot be included when this option is used.

Notes

To create a schema, the invoking user must have the CREATE privilege for the current database. (Of course, superusers bypass this check.)

Examples

Create a schema:

CREATE SCHEMA myschema;

Create a schema for user joe; the schema will also be named joe:

CREATE SCHEMA AUTHORIZATION joe;

Create a schema named test that will be owned by user joe, unless there already is a schema named test. (It does not matter whether joe owns the pre-existing schema.)

CREATE SCHEMA IF NOT EXISTS test AUTHORIZATION joe;

Create a schema and create a table and view within it:

CREATE SCHEMA hollywood
    CREATE TABLE films (title text, release date, awards text[])
    CREATE VIEW winners AS
        SELECT title, release FROM films WHERE awards IS NOT NULL;

Notice that the individual subcommands do not end with semicolons.

The following is an equivalent way of accomplishing the same result:

CREATE SCHEMA hollywood;
CREATE TABLE hollywood.films (title text, release date, awards text[]);
CREATE VIEW hollywood.winners AS
    SELECT title, release FROM hollywood.films WHERE awards IS NOT NULL;

Compatibility

The SQL standard allows a DEFAULT CHARACTER SET clause in CREATE SCHEMA, as well as more subcommand types than are presently accepted by PostgreSQL.

The SQL standard specifies that the subcommands in CREATE SCHEMA can appear in any order. The present PostgreSQL implementation does not handle all cases of forward references in subcommands; it might sometimes be necessary to reorder the subcommands in order to avoid forward references.
According to the SQL standard, the owner of a schema always owns all objects within it. PostgreSQL allows schemas to contain objects owned by users other than the schema owner. This can happen only if the schema owner grants the CREATE privilege on their schema to someone else, or a superuser chooses to create objects in it.

The **IF NOT EXISTS** option is a PostgreSQL extension.

**See Also**

ALTER SCHEMA, DROP SCHEMA
CREATE SEQUENCE

CREATE SEQUENCE — define a new sequence generator

Synopsis

CREATE [ TEMPORARY | TEMP ] SEQUENCE [ IF NOT EXISTS ] name
[ AS data_type ]
[ INCREMENT [ BY ] increment ]
[ MINVALUE minvalue | NO MINVALUE ] [ MAXVALUE maxvalue | NO MAXVALUE ]
[ START [ WITH ] start ] [ CACHE cache ] [ [ NO ] CYCLE ]
[ OWNED BY { table_name.column_name | NONE } ]

Description

CREATE SEQUENCE creates a new sequence number generator. This involves creating and initial-
izing a new special single-row table with the name name. The generator will be owned by the user
issuing the command.

If a schema name is given then the sequence is created in the specified schema. Otherwise it is created
in the current schema. Temporary sequences exist in a special schema, so a schema name cannot be
given when creating a temporary sequence. The sequence name must be distinct from the name of any
other sequence, table, index, view, or foreign table in the same schema.

After a sequence is created, you use the functions nextval, currval, and setval to operate on
the sequence. These functions are documented in Section 9.16.

Although you cannot update a sequence directly, you can use a query like:

SELECT * FROM name;

to examine the parameters and current state of a sequence. In particular, the last_value field of
the sequence shows the last value allocated by any session. (Of course, this value might be obsolete
by the time it's printed, if other sessions are actively doing nextval calls.)

Parameters

TEMPORARY or TEMP

If specified, the sequence object is created only for this session, and is automatically dropped on
session exit. Existing permanent sequences with the same name are not visible (in this session)
while the temporary sequence exists, unless they are referenced with schema-qualified names.

IF NOT EXISTS

Do not throw an error if a relation with the same name already exists. A notice is issued in this
case. Note that there is no guarantee that the existing relation is anything like the sequence that
would have been created - it might not even be a sequence.

name

The name (optionally schema-qualified) of the sequence to be created.
**CREATE SEQUENCE**

`data_type`

The optional clause `AS data_type` specifies the data type of the sequence. Valid types are `smallint`, `integer`, and `bigint`. `bigint` is the default. The data type determines the default minimum and maximum values of the sequence.

`increment`

The optional clause `INCREMENT BY increment` specifies which value is added to the current sequence value to create a new value. A positive value will make an ascending sequence, a negative one a descending sequence. The default value is 1.

`minvalue`

`NO MINVALUE`

The optional clause `MINVALUE minvalue` determines the minimum value a sequence can generate. If this clause is not supplied or `NO MINVALUE` is specified, then defaults will be used. The default for an ascending sequence is 1. The default for a descending sequence is the minimum value of the data type.

`maxvalue`

`NO MAXVALUE`

The optional clause `MAXVALUE maxvalue` determines the maximum value for the sequence. If this clause is not supplied or `NO MAXVALUE` is specified, then default values will be used. The default for an ascending sequence is the maximum value of the data type. The default for a descending sequence is -1.

`start`

The optional clause `START WITH start` allows the sequence to begin anywhere. The default starting value is `minvalue` for ascending sequences and `maxvalue` for descending ones.

`cache`

The optional clause `CACHE cache` specifies how many sequence numbers are to be preallocated and stored in memory for faster access. The minimum value is 1 (only one value can be generated at a time, i.e., no cache), and this is also the default.

`CYCLE`

`NO CYCLE`

The `CYCLE` option allows the sequence to wrap around when the `maxvalue` or `minvalue` has been reached by an ascending or descending sequence respectively. If the limit is reached, the next number generated will be the `minvalue` or `maxvalue`, respectively.

If `NO CYCLE` is specified, any calls to `nextval` after the sequence has reached its maximum value will return an error. If neither `CYCLE` or `NO CYCLE` are specified, `NO CYCLE` is the default.

`OWNED BY table_name.column_name`

`OWNED BY NONE`

The `OWNED BY` option causes the sequence to be associated with a specific table column, such that if that column (or its whole table) is dropped, the sequence will be automatically dropped as well. The specified table must have the same owner and be in the same schema as the sequence. `OWNED BY NONE`, the default, specifies that there is no such association.

**Notes**

Use `DROP SEQUENCE` to remove a sequence.
CREATE SEQUENCE

Sequences are based on bigint arithmetic, so the range cannot exceed the range of an eight-byte integer (-2^63 to 2^63-1).

Because nextval and setval calls are never rolled back, sequence objects cannot be used if "gapless" assignment of sequence numbers is needed. It is possible to build gapless assignment by using exclusive locking of a table containing a counter; but this solution is much more expensive than sequence objects, especially if many transactions need sequence numbers concurrently.

Unexpected results might be obtained if a cache setting greater than one is used for a sequence object that will be used concurrently by multiple sessions. Each session will allocate and cache successive sequence values during one access to the sequence object and increase the sequence object's last_value accordingly. Then, the next cache-1 uses of nextval within that session simply return the preallocated values without touching the sequence object. So, any numbers allocated but not used within a session will be lost when that session ends, resulting in "holes" in the sequence.

Furthermore, although multiple sessions are guaranteed to allocate distinct sequence values, the values might be generated out of sequence when all the sessions are considered. For example, with a cache setting of 10, session A might reserve values 1..10 and return nextval=1, then session B might reserve values 11..20 and return nextval=11 before session A has generated nextval=2. Thus, with a cache setting of one it is safe to assume that nextval values are generated sequentially; with a cache setting greater than one you should only assume that the nextval values are all distinct, not that they are generated purely sequentially. Also, last_value will reflect the latest value reserved by any session, whether or not it has yet been returned by nextval.

Another consideration is that a setval executed on such a sequence will not be noticed by other sessions until they have used up any preallocated values they have cached.

Examples

Create an ascending sequence called serial, starting at 101:

```
CREATE SEQUENCE serial START 101;
```

Select the next number from this sequence:

```
SELECT nextval('serial');
```

```
nextval
-------
101
```

Select the next number from this sequence:

```
SELECT nextval('serial');
```

```
nextval
-------
102
```

Use this sequence in an INSERT command:

```
INSERT INTO distributors VALUES (nextval('serial'), 'nothing');
```

Update the sequence value after a COPY FROM:
BEGIN;
COPY distributors FROM 'input_file';
SELECT setval('serial', max(id)) FROM distributors;
END;

Compatibility

CREATE SEQUENCE conforms to the SQL standard, with the following exceptions:

• Obtaining the next value is done using the nextval() function instead of the standard's NEXT VALUE FOR expression.

• The OWNED BY clause is a PostgreSQL extension.

See Also
ALTER SEQUENCE, DROP SEQUENCE
CREATE SERVER

CREATE SERVER — define a new foreign server

Synopsis

```
CREATE SERVER [ IF NOT EXISTS ] server_name [ TYPE 'server_type' ]
[ VERSION 'server_version' ]
FOREIGN DATA WRAPPER fdw_name
[ OPTIONS ( option 'value' [, ... ] ) ]
```

Description

CREATE SERVER defines a new foreign server. The user who defines the server becomes its owner.

A foreign server typically encapsulates connection information that a foreign-data wrapper uses to access an external data resource. Additional user-specific connection information may be specified by means of user mappings.

The server name must be unique within the database.

Creating a server requires USAGE privilege on the foreign-data wrapper being used.

Parameters

**IF NOT EXISTS**

Do not throw an error if a server with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing server is anything like the one that would have been created.

**server_name**

The name of the foreign server to be created.

**server_type**

Optional server type, potentially useful to foreign-data wrappers.

**server_version**

Optional server version, potentially useful to foreign-data wrappers.

**fdw_name**

The name of the foreign-data wrapper that manages the server.

**OPTIONS ( option 'value' [, ... ] )**

This clause specifies the options for the server. The options typically define the connection details of the server, but the actual names and values are dependent on the server's foreign-data wrapper.

Notes

When using the dblink module, a foreign server's name can be used as an argument of the dblink_connect function to indicate the connection parameters. It is necessary to have the USAGE privilege on the foreign server to be able to use it in this way.
CREATE SERVER

Examples

Create a server myserver that uses the foreign-data wrapper postgres_fdw:

```
CREATE SERVER myserver FOREIGN DATA WRAPPER postgres_fdw OPTIONS
    (host 'foo', dbname 'foodb', port '5432');
```

See postgres_fdw for more details.

Compatibility

CREATE SERVER conforms to ISO/IEC 9075-9 (SQL/MED).

See Also

ALTER SERVER, DROP SERVER, CREATE FOREIGN DATA WRAPPER, CREATE FOREIGN TABLE, CREATE USER MAPPING
CREATE STATISTICS

CREATE STATISTICS — define extended statistics

Synopsis

```
CREATE STATISTICS [ IF NOT EXISTS ] statistics_name
   [ ( statistics_kind [, ... ] ) ]
ON column_name, column_name [, ...]
FROM table_name
```

Description

CREATE STATISTICS will create a new extended statistics object tracking data about the specified table, foreign table or materialized view. The statistics object will be created in the current database and will be owned by the user issuing the command.

If a schema name is given (for example, `CREATE STATISTICS myschema.mystat ...`) then the statistics object is created in the specified schema. Otherwise it is created in the current schema. The name of the statistics object must be distinct from the name of any other statistics object in the same schema.

Parameters

**IF NOT EXISTS**

Do not throw an error if a statistics object with the same name already exists. A notice is issued in this case. Note that only the name of the statistics object is considered here, not the details of its definition.

**statistics_name**

The name (optionally schema-qualified) of the statistics object to be created.

**statistics_kind**

A statistics kind to be computed in this statistics object. Currently supported kinds are `ndistinct`, which enables n-distinct statistics, and `dependencies`, which enables functional dependency statistics. If this clause is omitted, all supported statistics kinds are included in the statistics object. For more information, see Section 14.2.2 and Section 70.2.

**column_name**

The name of a table column to be covered by the computed statistics. At least two column names must be given.

**table_name**

The name (optionally schema-qualified) of the table containing the column(s) the statistics are computed on.

Notes

You must be the owner of a table to create a statistics object reading it. Once created, however, the ownership of the statistics object is independent of the underlying table(s).
CREATE STATISTICS

Examples

Create table `t1` with two functionally dependent columns, i.e. knowledge of a value in the first column is sufficient for determining the value in the other column. Then functional dependency statistics are built on those columns:

```sql
CREATE TABLE t1 (
    a   int,
    b   int
);

INSERT INTO t1 SELECT i/100, i/500
    FROM generate_series(1,1000000) s(i);

ANALYZE t1;

-- the number of matching rows will be drastically underestimated:
EXPLAIN ANALYZE SELECT * FROM t1 WHERE (a = 1) AND (b = 0);

CREATE STATISTICS s1 (dependencies) ON a, b FROM t1;

ANALYZE t1;

-- now the row count estimate is more accurate:
EXPLAIN ANALYZE SELECT * FROM t1 WHERE (a = 1) AND (b = 0);
```

Without functional-dependency statistics, the planner would assume that the two `WHERE` conditions are independent, and would multiply their selectivities together to arrive at a much-too-small row count estimate. With such statistics, the planner recognizes that the `WHERE` conditions are redundant and does not underestimate the row count.

Compatibility

There is no `CREATE STATISTICS` command in the SQL standard.

See Also

`ALTER STATISTICS`, `DROP STATISTICS`
CREATE SUBSCRIPTION

CREATE SUBSCRIPTION — define a new subscription

Synopsis

CREATE SUBSCRIPTION subscription_name
  CONNECTION 'conninfo'
  PUBLICATION publication_name [, ...]
  [ WITH ( subscription_parameter [= value] [, ... ] ) ]

Description

CREATE SUBSCRIPTION adds a new subscription for the current database. The subscription name
must be distinct from the name of any existing subscription in the database.

The subscription represents a replication connection to the publisher. As such this command does not
only add definitions in the local catalogs but also creates a replication slot on the publisher.

A logical replication worker will be started to replicate data for the new subscription at the commit
of the transaction where this command is run.

Additional information about subscriptions and logical replication as a whole is available at Sec-
tion 31.2 and Chapter 31.

Parameters

subscription_name
  The name of the new subscription.

CONNECTION 'conninfo'
  The connection string to the publisher. For details see Section 34.1.1.

PUBLICATION publication_name
  Names of the publications on the publisher to subscribe to.

WITH ( subscription_parameter [= value] [, ... ] )
  This clause specifies optional parameters for a subscription. The following parameters are sup-
  ported:

  copy_data (boolean)
    Specifies whether the existing data in the publications that are being subscribed to should be
copied once the replication starts. The default is true.

  create_slot (boolean)
    Specifies whether the command should create the replication slot on the publisher. The de-
    fault is true.

  enabled (boolean)
    Specifies whether the subscription should be actively replicating, or whether it should be just
    setup but not started yet. The default is true.
CREATE SUBSCRIPTION

slot_name (string)

Name of the replication slot to use. The default behavior is to use the name of the subscription for the slot name.

When slot_name is set to NONE, there will be no replication slot associated with the subscription. This can be used if the replication slot will be created later manually. Such subscriptions must also have both enabled and create_slot set to false.

synchronous_commit (enum)

The value of this parameter overrides the synchronous_commit setting. The default value is off.

It is safe to use off for logical replication: If the subscriber loses transactions because of missing synchronization, the data will be resent from the publisher.

A different setting might be appropriate when doing synchronous logical replication. The logical replication workers report the positions of writes and flushes to the publisher, and when using synchronous replication, the publisher will wait for the actual flush. This means that setting synchronous_commit for the subscriber to off when the subscription is used for synchronous replication might increase the latency for COMMIT on the publisher. In this scenario, it can be advantageous to set synchronous_commit to local or higher.

connect (boolean)

Specifies whether the CREATE SUBSCRIPTION should connect to the publisher at all. Setting this to false will change default values of enabled, create_slot and copy_data to false.

It is not allowed to combine connect set to false and enabled, create_slot, or copy_data set to true.

Since no connection is made when this option is set to false, the tables are not subscribed, and so after you enable the subscription nothing will be replicated. It is required to run ALTER SUBSCRIPTION ... REFRESH PUBLICATION in order for tables to be subscribed.

Notes

See Section 31.7 for details on how to configure access control between the subscription and the publication instance.

When creating a replication slot (the default behavior), CREATE SUBSCRIPTION cannot be executed inside a transaction block.

Creating a subscription that connects to the same database cluster (for example, to replicate between databases in the same cluster or to replicate within the same database) will only succeed if the replication slot is not created as part of the same command. Otherwise, the CREATE SUBSCRIPTION call will hang. To make this work, create the replication slot separately (using the function pg_create_logical_replication_slot with the plugin name pgoutput) and create the subscription using the parameter create_slot = false. This is an implementation restriction that might be lifted in a future release.

Examples

Create a subscription to a remote server that replicates tables in the publications mypublication and insert_only and starts replicating immediately on commit:

CREATE SUBSCRIPTION mysub
CREATE SUBSCRIPTION
CONNECTION 'host=192.168.1.50 port=5432 user=foo
dbname=foodb'
PUBLICACION mypublication, insert_only;

Create a subscription to a remote server that replicates tables in the insert_only publication and does not start replicating until enabled at a later time.

CREATE SUBSCRIPTION mysub
CONNECTION 'host=192.168.1.50 port=5432 user=foo
dbname=foodb'
PUBLICACION insert_only
WITH (enabled = false);

Compatibility

CREATE SUBSCRIPTION is a PostgreSQL extension.

See Also

ALTER SUBSCRIPTION, DROP SUBSCRIPTION, CREATE PUBLICATION, ALTER PUBLICATION
CREATE TABLE

CREATE TABLE — define a new table

Synopsis

```
CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE
[ IF NOT EXISTS ] table_name ( [ column_name data_type [ COLLATE collation ] [ column_constraint [ ... ] ] ]
  [ table_constraint ]
  [ LIKE source_table [ like_option ... ] ]
  [ , ... ] )
  [ INHERITS ( parent_table [ , ... ] ) ]
[ PARTITION BY { RANGE | LIST | HASH } ( { column_name | ( expression ) ) [ COLLATE collation ] [ opclass ] [, ... ] } ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]
```

```
CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE
[ IF NOT EXISTS ] table_name
  OF type_name ( [ column_name [ WITH OPTIONS ] [ column_constraint [ ... ] ] ]
  [ table_constraint ]
  [ , ... ] )
[ PARTITION BY { RANGE | LIST | HASH } ( { column_name | ( expression ) ) [ COLLATE collation ] [ opclass ] [, ... ] } ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]
```

```
CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE
[ IF NOT EXISTS ] table_name
  PARTITION OF parent_table ( [ column_name [ WITH OPTIONS ] [ column_constraint [ ... ] ] ]
  [ table_constraint ]
  [ , ... ] )
[ FOR VALUES partition_bound_spec | DEFAULT ]
[ PARTITION BY { RANGE | LIST | HASH } ( { column_name | ( expression ) ) [ COLLATE collation ] [ opclass ] [, ... ] } ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]
```

where column_constraint is:

```
[ CONSTRAINT constraint_name ]
[ NOT NULL ]
NULL |
```

1581
CREATE TABLE

CHECK ( expression ) [ NO INHERIT ] |
DEFAULT default_expr |
GENERATED { ALWAYS | BY DEFAULT } AS IDENTITY |
[ { sequence_options } ] |
UNIQUE index_parameters |
PRIMARY KEY index_parameters |
REFERENCES reftable [ ( refcolumn ) ] [ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ] |
[ ON DELETE action ] [ ON UPDATE action ] |
[ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ] |
and table_constraint is:

[ CONSTRAINT constraint_name ]
CHECK ( expression ) [ NO INHERIT ] |
UNIQUE ( column_name [, ... ] ) index_parameters |
PRIMARY KEY ( column_name [, ... ] ) index_parameters |
EXCLUDE [ USING index_method ] ( exclude_element WITH operator [, ... ] ) index_parameters [ WHERE ( predicate ) ] |
FOREIGN KEY ( column_name [, ... ] ) REFERENCES reftable |
[ ( refcolumn [, ... ] ) ] |
[ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ] [ ON DELETE action ] [ ON UPDATE action ] |
[ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ] |
and like_option is:

{ INCLUDING | EXCLUDING } { COMMENTS | CONSTRAINTS | DEFAULTS | IDENTITY | INDEXES | STATISTICS | STORAGE | ALL }
and partition_bound_spec is:

IN ( ( numeric_literal | string_literal | TRUE | FALSE | NULL ) [, ... ] ) |
FROM ( ( numeric_literal | string_literal | TRUE | FALSE | MINVALUE | MAXVALUE ) [, ... ] ) |
TO ( ( numeric_literal | string_literal | TRUE | FALSE | MINVALUE | MAXVALUE ) [, ... ] ) |
WITH ( MODULUS numeric_literal, REMAINDER numeric_literal )

index_parameters in UNIQUE, PRIMARY KEY, and EXCLUDE constraints are:

[ INCLUDE ( column_name [, ... ] ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ USING INDEX TABLESPACE tablespace_name ]

exclude_element in an EXCLUDE constraint is:

{ column_name | ( expression ) } [ opclass ] [ ASC | DESC ] [ NULLS ( FIRST | LAST ) ]

Description

CREATE TABLE will create a new, initially empty table in the current database. The table will be owned by the user issuing the command.
CREATE TABLE

If a schema name is given (for example, CREATE TABLE myschema.mytable ...) then the table is created in the specified schema. Otherwise it is created in the current schema. Temporary tables exist in a special schema, so a schema name cannot be given when creating a temporary table. The name of the table must be distinct from the name of any other table, sequence, index, view, or foreign table in the same schema.

CREATE TABLE also automatically creates a data type that represents the composite type corresponding to one row of the table. Therefore, tables cannot have the same name as any existing data type in the same schema.

The optional constraint clauses specify constraints (tests) that new or updated rows must satisfy for an insert or update operation to succeed. A constraint is an SQL object that helps define the set of valid values in the table in various ways.

There are two ways to define constraints: table constraints and column constraints. A column constraint is defined as part of a column definition. A table constraint definition is not tied to a particular column, and it can encompass more than one column. Every column constraint can also be written as a table constraint; a column constraint is only a notational convenience for use when the constraint only affects one column.

To be able to create a table, you must have USAGE privilege on all column types or the type in the OF clause, respectively.

Parameters

TEMPORARY or TEMP

If specified, the table is created as a temporary table. Temporary tables are automatically dropped at the end of a session, or optionally at the end of the current transaction (see ON COMMIT below). Existing permanent tables with the same name are not visible to the current session while the temporary table exists, unless they are referenced with schema-qualified names. Any indexes created on a temporary table are automatically temporary as well.

The autovacuum daemon cannot access and therefore cannot vacuum or analyze temporary tables. For this reason, appropriate vacuum and analyze operations should be performed via session SQL commands. For example, if a temporary table is going to be used in complex queries, it is wise to run ANALYZE on the temporary table after it is populated.

Optionally, GLOBAL or LOCAL can be written before TEMPORARY or TEMP. This presently makes no difference in PostgreSQL and is deprecated; see Compatibility.

UNLOGGED

If specified, the table is created as an unlogged table. Data written to unlogged tables is not written to the write-ahead log (see Chapter 30), which makes them considerably faster than ordinary tables. However, they are not crash-safe: an unlogged table is automatically truncated after a crash or unclean shutdown. The contents of an unlogged table are also not replicated to standby servers. Any indexes created on an unlogged table are automatically unlogged as well.

IF NOT EXISTS

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing relation is anything like the one that would have been created.

table_name

The name (optionally schema-qualified) of the table to be created.
CREATE TABLE

OF type_name

Creates a typed table, which takes its structure from the specified composite type (name optionally
schema-qualified). A typed table is tied to its type; for example the table will be dropped if the
type is dropped (with DROP TYPE ... CASCADE).

When a typed table is created, then the data types of the columns are determined by the underlying
composite type and are not specified by the CREATE TABLE command. But the CREATE TABLE
command can add defaults and constraints to the table and can specify storage parameters.

column_name

The name of a column to be created in the new table.

data_type

The data type of the column. This can include array specifiers. For more information on the data
types supported by PostgreSQL, refer to Chapter 8.

COLLATE collation

The COLLATE clause assigns a collation to the column (which must be of a collatable data type).
If not specified, the column data type's default collation is used.

INHERITS ( parent_table [, ... ] )

The optional INHERITS clause specifies a list of tables from which the new table automatically
inherits all columns. Parent tables can be plain tables or foreign tables.

Use of INHERITS creates a persistent relationship between the new child table and its parent
table(s). Schema modifications to the parent(s) normally propagate to children as well, and by
default the data of the child table is included in scans of the parent(s).

If the same column name exists in more than one parent table, an error is reported unless the data
types of the columns match in each of the parent tables. If there is no conflict, then the duplicate
columns are merged to form a single column in the new table. If the column name list of the new
table contains a column name that is also inherited, the data type must likewise match the inherited
column(s), and the column definitions are merged into one. If the new table explicitly specifies
a default value for the column, this default overrides any defaults from inherited declarations of
the column. Otherwise, any parents that specify default values for the column must all specify the
same default, or an error will be reported.

CHECK constraints are merged in essentially the same way as columns: if multiple parent tables
and/or the new table definition contain identically-named CHECK constraints, these constraints
must all have the same check expression, or an error will be reported. Constraints having the
same name and expression will be merged into one copy. A constraint marked NO INHERIT in
a parent will not be considered. Notice that an unnamed CHECK constraint in the new table will
never be merged, since a unique name will always be chosen for it.

Column STORAGE settings are also copied from parent tables.

If a column in the parent table is an identity column, that property is not inherited. A column in
the child table can be declared identity column if desired.

PARTITION BY { RANGE | LIST | HASH } ( { column_name | ( expression ) } [ opclass ] [, ... ] )

The optional PARTITION BY clause specifies a strategy of partitioning the table. The table
thus created is called a partitioned table. The parenthesized list of columns or expressions forms
the partition key for the table. When using range or hash partitioning, the partition key can in-
clude multiple columns or expressions (up to 32, but this limit can be altered when building Post-
greSQL), but for list partitioning, the partition key must consist of a single column or expression.
CREATE TABLE

Range and list partitioning require a btree operator class, while hash partitioning requires a hash operator class. If no operator class is specified explicitly, the default operator class of the appropriate type will be used; if no default operator class exists, an error will be raised. When hash partitioning is used, the operator class used must implement support function 2 (see Section 38.15.3 for details).

A partitioned table is divided into sub-tables (called partitions), which are created using separate CREATE TABLE commands. The partitioned table is itself empty. A data row inserted into the table is routed to a partition based on the value of columns or expressions in the partition key. If no existing partition matches the values in the new row, an error will be reported.

Partitioned tables do not support EXCLUDE constraints; however, you can define these constraints on individual partitions. Also, while it's possible to define PRIMARY KEY constraints on partitioned tables, creating foreign keys that reference a partitioned table is not yet supported.

See Section 5.10 for more discussion on table partitioning.

PARTITION OF parent_table { FOR VALUES partition_bound_spec | DEFAULT }

Creates the table as a partition of the specified parent table. The table can be created either as a partition for specific values using FOR VALUES or as a default partition using DEFAULT. This option is not available for hash-partitioned tables.

The partition_bound_spec must correspond to the partitioning method and partition key of the parent table, and must not overlap with any existing partition of that parent. The form with IN is used for list partitioning, the form with FROM and TO is used for range partitioning, and the form with WITH is used for hash partitioning.

Each of the values specified in the partition_bound_spec is a literal, NULL, MINVALUE, or MAXVALUE. Each literal value must be either a numeric constant that is coercible to the corresponding partition key column's type, or a string literal that is valid input for that type.

When creating a list partition, NULL can be specified to signify that the partition allows the partition key column to be null. However, there cannot be more than one such list partition for a given parent table. NULL cannot be specified for range partitions.

When creating a range partition, the lower bound specified with FROM is an inclusive bound, whereas the upper bound specified with TO is an exclusive bound. That is, the values specified in the FROM list are valid values of the corresponding partition key columns for this partition, whereas those in the TO list are not. Note that this statement must be understood according to the rules of row-wise comparison (Section 9.23.5). For example, given PARTITION BY RANGE (x, y), a partition bound FROM (1, 2) TO (3, 4) allows x=1 with any y>=2, x=2 with any non-null y, and x=3 with any y<4.

The special values MINVALUE and MAXVALUE may be used when creating a range partition to indicate that there is no lower or upper bound on the column's value. For example, a partition defined using FROM (MINVALUE) TO (10) allows any values less than 10, and a partition defined using FROM (10) TO (MAXVALUE) allows any values greater than or equal to 10.

When creating a range partition involving more than one column, it can also make sense to use MAXVALUE as part of the lower bound, and MINVALUE as part of the upper bound. For example, a partition defined using FROM (0, MAXVALUE) TO (10, MAXVALUE) allows any rows where the first partition key column is greater than 0 and less than or equal to 10. Similarly, a partition defined using FROM ('a', MINVALUE) TO ('b', MINVALUE) allows any rows where the first partition key column starts with "a".

Note that if MINVALUE or MAXVALUE is used for one column of a partitioning bound, the same value must be used for all subsequent columns. For example, (10, MINVALUE, 0) is not a valid bound; you should write (10, MINVALUE, MINVALUE).
CREATE TABLE

Also note that some element types, such as timestamp, have a notion of "infinity", which is just
another value that can be stored. This is different from MINVALUE and MAXVALUE, which are
not real values that can be stored, but rather they are ways of saying that the value is unbounded.
MAXVALUE can be thought of as being greater than any other value, including "infinity" and
MINVALUE as being less than any other value, including "minus infinity". Thus the range FROM
('infinity') TO (MAXVALUE) is not an empty range; it allows precisely one value to be stored — "infinity".

If DEFAULT is specified, the table will be created as a default partition of the parent table. The
parent can either be a list or range partitioned table. A partition key value not fitting into any other
partition of the given parent will be routed to the default partition. There can be only one default
partition for a given parent table.

When a table has an existing DEFAULT partition and a new partition is added to it, the existing
default partition must be scanned to verify that it does not contain any rows which properly belong
in the new partition. If the default partition contains a large number of rows, this may be slow.
The scan will be skipped if the default partition is a foreign table or if it has a constraint which
proves that it cannot contain rows which should be placed in the new partition.

When creating a hash partition, a modulus and remainder must be specified. The modulus must
be a positive integer, and the remainder must be a non-negative integer less than the modulus.
Typically, when initially setting up a hash-partitioned table, you should choose a modulus equal
to the number of partitions and assign every table the same modulus and a different remainder
(see examples, below). However, it is not required that every partition have the same modulus,
only that every modulus which occurs among the partitions of a hash-partitioned table is a factor
of the next larger modulus. This allows the number of partitions to be increased incrementally
without needing to move all the data at once. For example, suppose you have a hash-partitioned
table with 8 partitions, each of which has modulus 8, but find it necessary to increase the number
of partitions to 16. You can detach one of the modulus-8 partitions, create two new modulus-16
partitions covering the same portion of the key space (one with a remainder equal to the remainder
of the detached partition, and the other with a remainder equal to that value plus 8), and repopulate
them with data. You can then repeat this -- perhaps at a later time -- for each modulus-8 partition
until none remain. While this may still involve a large amount of data movement at each step, it
is still better than having to create a whole new table and move all the data at once.

A partition must have the same column names and types as the partitioned table to which it be-
longs. If the parent is specified WITH OIDS then all partitions must have OIDs; the parent's
OID column will be inherited by all partitions just like any other column. Modifications to the
column names or types of a partitioned table, or the addition or removal of an OID column, will
automatically propagate to all partitions. CHECK constraints will be inherited automatically by
every partition, but an individual partition may specify additional CHECK constraints; additional
constraints with the same name and condition as in the parent will be merged with the parent
constraint. Defaults may be specified separately for each partition.

Rows inserted into a partitioned table will be automatically routed to the correct partition. If no
suitable partition exists, an error will occur.

Operations such as TRUNCATE which normally affect a table and all of its inheritance children
will cascade to all partitions, but may also be performed on an individual partition. Note that
dropping a partition with DROP TABLE requires taking an ACCESS EXCLUSIVE lock on the
parent table.

LIKE source_table [ like_option ... ]

The LIKE clause specifies a table from which the new table automatically copies all column
names, their data types, and their not-null constraints.

Unlike INHERITS, the new table and original table are completely decoupled after creation is
complete. Changes to the original table will not be applied to the new table, and it is not possible
to include data of the new table in scans of the original table.
CREATE TABLE

Default expressions for the copied column definitions will be copied only if **INCLUDING DEFAULTS** is specified. The default behavior is to exclude default expressions, resulting in the copied columns in the new table having null defaults. Note that copying defaults that call database-modification functions, such as `nextval`, may create a functional linkage between the original and new tables.

Any identity specifications of copied column definitions will only be copied if **INCLUDING IDENTITY** is specified. A new sequence is created for each identity column of the new table, separate from the sequences associated with the old table.

Not-null constraints are always copied to the new table. **CHECK** constraints will be copied only if **INCLUDING CONSTRAINTS** is specified. No distinction is made between column constraints and table constraints.

Extended statistics are copied to the new table if **INCLUDING STATISTICS** is specified.

Indexes, **PRIMARY KEY**, **UNIQUE**, and **EXCLUDE** constraints on the original table will be created on the new table only if **INCLUDING INDEXES** is specified. Names for the new indexes and constraints are chosen according to the default rules, regardless of how the originals were named. (This behavior avoids possible duplicate-name failures for the new indexes.)

**STORAGE** settings for the copied column definitions will be copied only if **INCLUDING STORAGE** is specified. The default behavior is to exclude **STORAGE** settings, resulting in the copied columns in the new table having type-specific default settings. For more on **STORAGE** settings, see Section 68.2.

Comments for the copied columns, constraints, and indexes will be copied only if **INCLUDING COMMENTS** is specified. The default behavior is to exclude comments, resulting in the copied columns and constraints in the new table having no comments.

**INCLUDING ALL** is an abbreviated form of **INCLUDING COMMENTS INCLUDING CONSTRAINTS INCLUDING DEFAULTS INCLUDING IDENTITY INCLUDING INDEXES INCLUDING STATISTICS INCLUDING STORAGE**.

Note that unlike **INHERITS**, columns and constraints copied by **LIKE** are not merged with similarly named columns and constraints. If the same name is specified explicitly or in another **LIKE** clause, an error is signaled.

The **LIKE** clause can also be used to copy column definitions from views, foreign tables, or composite types. Inapplicable options (e.g., **INCLUDING INDEXES** from a view) are ignored.

**CONSTRAINT constraint_name**

An optional name for a column or table constraint. If the constraint is violated, the constraint name is present in error messages, so constraint names like `col must be positive` can be used to communicate helpful constraint information to client applications. (Double-quotes are needed to specify constraint names that contain spaces.) If a constraint name is not specified, the system generates a name.

**NOT NULL**

The column is not allowed to contain null values.

**NULL**

The column is allowed to contain null values. This is the default.

This clause is only provided for compatibility with non-standard SQL databases. Its use is discouraged in new applications.

**CHECK ( expression ) [ NO INHERIT ]**

The **CHECK** clause specifies an expression producing a Boolean result which new or updated rows must satisfy for an insert or update operation to succeed. Expressions evaluating to `TRUE` or
UNKNOWN succeed. Should any row of an insert or update operation produce a FALSE result, an error exception is raised and the insert or update does not alter the database. A check constraint specified as a column constraint should reference that column's value only, while an expression appearing in a table constraint can reference multiple columns.

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than columns of the current row. The system column `tableoid` may be referenced, but not any other system column.

A constraint marked with `NO INHERIT` will not propagate to child tables.

When a table has multiple CHECK constraints, they will be tested for each row in alphabetical order by name, after checking NOT NULL constraints. (PostgreSQL versions before 9.5 did not honor any particular firing order for CHECK constraints.)

**DEFAULT** `default_expr`

The DEFAULT clause assigns a default data value for the column whose column definition it appears within. The value is any variable-free expression (subqueries and cross-references to other columns in the current table are not allowed). The data type of the default expression must match the data type of the column.

The default expression will be used in any insert operation that does not specify a value for the column. If there is no default for a column, then the default is null.

**GENERATED { ALWAYS | BY DEFAULT } AS IDENTITY [ ( sequence_options ) ]**

This clause creates the column as an identity column. It will have an implicit sequence attached to it and the column in new rows will automatically have values from the sequence assigned to it.

The clauses ALWAYS and BY DEFAULT determine how the sequence value is given precedence over a user-specified value in an INSERT statement. If ALWAYS is specified, a user-specified value is only accepted if the INSERT statement specifies OVERRIDING SYSTEM VALUE. If BY DEFAULT is specified, then the user-specified value takes precedence. See INSERT for details. (In the COPY command, user-specified values are always used regardless of this setting.)

The optional sequence_options clause can be used to override the options of the sequence. See CREATE SEQUENCE for details.

**UNIQUE (column constraint)**

**UNIQUE ( column_name [, ... ] ) [ INCLUDE ( column_name [, ...] ) ] (table constraint)**

The UNIQUE constraint specifies that a group of one or more columns of a table can contain only unique values. The behavior of the unique table constraint is the same as that for column constraints, with the additional capability to span multiple columns.

For the purpose of a unique constraint, null values are not considered equal.

Each unique table constraint must name a set of columns that is different from the set of columns named by any other unique or primary key constraint defined for the table. (Otherwise it would just be the same constraint listed twice.)

When establishing a unique constraint for a multi-level partition hierarchy, all the columns in the partition key of the target partitioned table, as well as those of all its descendant partitioned tables, must be included in the constraint definition.

Adding a unique constraint will automatically create a unique btree index on the column or group of columns used in the constraint. The optional clause INCLUDE adds to that index one or more columns on which the uniqueness is not enforced. Note that although the constraint is not enforced on the included columns, it still depends on them. Consequently, some operations on these columns (e.g. DROP COLUMN) can cause cascaded constraint and index deletion.
The \texttt{CREATE TABLE} command is used to create a new table in a database. Each table must have at least one primary key constraint, \texttt{PRIMARY KEY}, which specifies the columns that uniquely identify each row in the table.

The \texttt{PRIMARY KEY} constraint specifies that a column or columns of a table can contain only unique (non-duplicate), nonnull values. Only one primary key can be specified for a table, whether as a column constraint or a table constraint.

The primary key constraint should name a set of columns that is different from the set of columns named by any unique constraint defined for the same table. (Otherwise, the unique constraint is redundant and will be discarded.)

\texttt{PRIMARY KEY} enforces the same data constraints as a combination of \texttt{UNIQUE} and \texttt{NOT NULL}, but identifying a set of columns as the primary key also provides metadata about the design of the schema, since a primary key implies that other tables can rely on this set of columns as a unique identifier for rows.

\texttt{PRIMARY KEY} constraints share the restrictions that \texttt{UNIQUE} constraints have when placed on partitioned tables.

Adding a \texttt{PRIMARY KEY} constraint will automatically create a unique btree index on the column or group of columns used in the constraint. The optional \texttt{INCLUDE} clause allows a list of columns to be specified which will be included in the non-key portion of the index. Although uniqueness is not enforced on the included columns, the constraint still depends on them. Consequently, some operations on the included columns (e.g. \texttt{DROP COLUMN}) can cause cascaded constraint and index deletion.

Exclusion constraints are implemented using an index, so each specified operator must be associated with an appropriate operator class (see Section 11.9) for the index access method \texttt{index_method}. The operators are required to be commutative. Each \texttt{exclude_element} can optionally specify an operator class and/or ordering options; these are described fully under \texttt{CREATE INDEX}.

Exclusion constraints are implemented using an index, so each specified operator must be associated with an appropriate operator class (see Section 11.9) for the index access method \texttt{index_method}. The operators are required to be commutative. Each \texttt{exclude_element} can optionally specify an operator class and/or ordering options; these are described fully under \texttt{CREATE INDEX}.

The access method must support \texttt{amgettuple} (see Chapter 61); at present this means GIN cannot be used. Although it's allowed, there is little point in using B-tree or hash indexes with an exclusion constraint, because this does nothing that an ordinary unique constraint doesn't do better. So in practice the access method will always be GiST or SP-GiST.

The \texttt{predicate} allows you to specify an exclusion constraint on a subset of the table; internally this creates a partial index. Note that parentheses are required around the predicate.

These clauses specify a foreign key constraint, which requires that a group of one or more columns of the new table must only contain values that match values in the referenced column(s) of some
row of the referenced table. If the refcolumn list is omitted, the primary key of the reftable is used. The referenced columns must be the columns of a non-deferrable unique or primary key constraint in the referenced table. The user must have REFERENCES permission on the referenced table (either the whole table, or the specific referenced columns). The addition of a foreign key constraint requires a SHARE ROW EXCLUSIVE lock on the referenced table. Note that foreign key constraints cannot be defined between temporary tables and permanent tables. Also note that while it is possible to define a foreign key on a partitioned table, it is not possible to declare a foreign key that references a partitioned table.

A value inserted into the referencing column(s) is matched against the values of the referenced table and referenced columns using the given match type. There are three match types: MATCH FULL, MATCH PARTIAL, and MATCH SIMPLE (which is the default). MATCH FULL will not allow one column of a multicolumn foreign key to be null unless all foreign key columns are null; if they are all null, the row is not required to have a match in the referenced table. MATCH SIMPLE allows any of the foreign key columns to be null; if any of them are null, the row is not required to have a match in the referenced table. MATCH PARTIAL is not yet implemented. (Of course, NOT NULL constraints can be applied to the referencing column(s) to prevent these cases from arising.)

In addition, when the data in the referenced columns is changed, certain actions are performed on the data in this table's columns. The ON DELETE clause specifies the action to perform when a referenced row in the referenced table is being deleted. Likewise, the ON UPDATE clause specifies the action to perform when a referenced column in the referenced table is being updated to a new value. If the row is updated, but the referenced column is not actually changed, no action is done. Referential actions other than the NO ACTION check cannot be deferred, even if the constraint is declared deferrable. There are the following possible actions for each clause:

**NO ACTION**

Produce an error indicating that the deletion or update would create a foreign key constraint violation. If the constraint is deferred, this error will be produced at constraint check time if there still exist any referencing rows. This is the default action.

**RESTRICT**

Produce an error indicating that the deletion or update would create a foreign key constraint violation. This is the same as NO ACTION except that the check is not deferrable.

**CASCADE**

Delete any rows referencing the deleted row, or update the values of the referencing column(s) to the new values of the referenced columns, respectively.

**SET NULL**

Set the referencing column(s) to null.

**SET DEFAULT**

Set the referencing column(s) to their default values. (There must be a row in the referenced table matching the default values, if they are not null, or the operation will fail.)

If the referenced column(s) are changed frequently, it might be wise to add an index to the referencing column(s) so that referential actions associated with the foreign key constraint can be performed more efficiently.

**DEFERRABLE**

**NOT DEFERRABLE**

This controls whether the constraint can be deferred. A constraint that is not deferrable will be checked immediately after every command. Checking of constraints that are deferrable can be
CREATE TABLE postponed until the end of the transaction (using the SET CONSTRAINTS command). NOT DE-
FERRABLE is the default. Currently, only UNIQUE, PRIMARY KEY, EXCLUDE, and REFER-
ENCES (foreign key) constraints accept this clause. NOT NULL and CHECK constraints are not
deferrable. Note that deferrable constraints cannot be used as conflict arbitrators in an INSERT
statement that includes an ON CONFLICT DO UPDATE clause.

INITIALLY IMMEDIATE
INITIALLY DEFERRED

If a constraint is deferrable, this clause specifies the default time to check the constraint. If the
constraint is INITIALLY IMMEDIATE, it is checked after each statement. This is the default. If
the constraint is INITIALLY DEFERRED, it is checked only at the end of the transaction. The
constraint check time can be altered with the SET CONSTRAINTS command.

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for a table or index; see Storage Parameters for
more information. The WITH clause for a table can also include OIDS=TRUE (or just OIDS)
to specify that rows of the new table should have OIDs (object identifiers) assigned to them,
or OIDS=FALSE to specify that the rows should not have OIDs. If OIDS is not specified, the
default setting depends upon the default_with_oids configuration parameter. (If the new table
inherits from any tables that have OIDs, then OIDS=TRUE is forced even if the command says
OIDS=FALSE.)

If OIDS=FALSE is specified or implied, the new table does not store OIDs and no OID will be as-
signed for a row inserted into it. This is generally considered worthwhile, since it will reduce OID
consumption and thereby postpone the wraparound of the 32-bit OID counter. Once the counter
wraps around, OIDs can no longer be assumed to be unique, which makes them considerably less
useful. In addition, excluding OIDs from a table reduces the space required to store the table on
disk by 4 bytes per row (on most machines), slightly improving performance.

To remove OIDs from a table after it has been created, use ALTER TABLE.

WITH OIDS
WITHOUT OIDS

These are obsolescent syntaxes equivalent to WITH (OIDS) and WITH (OIDS=FALSE),
respectively. If you wish to give both an OIDS setting and storage parameters, you must use the
WITH ( ... ) syntax; see above.

ON COMMIT

The behavior of temporary tables at the end of a transaction block can be controlled using ON
COMMIT. The three options are:

PRESERVE ROWS

No special action is taken at the ends of transactions. This is the default behavior.

DELETE ROWS

All rows in the temporary table will be deleted at the end of each transaction block. Essential-
ly, an automatic TRUNCATE is done at each commit.

DROP

The temporary table will be dropped at the end of the current transaction block.

TABLESPACE tablespace_name

The tablespace_name is the name of the tablespace in which the new table is to be created.
If not specified, default_tablespace is consulted, or temp_tablespaces if the table is temporary.
CREATE TABLE

USING INDEX TABLESPACE tablespace_name

This clause allows selection of the tablespace in which the index associated with a UNIQUE, PRIMARY KEY, or EXCLUDE constraint will be created. If not specified, default_tablespace is consulted, or temp_tablespaces if the table is temporary.

Storage Parameters

The WITH clause can specify storage parameters for tables, and for indexes associated with a UNIQUE, PRIMARY KEY, or EXCLUDE constraint. Storage parameters for indexes are documented in CREATE INDEX. The storage parameters currently available for tables are listed below. For many of these parameters, as shown, there is an additional parameter with the same name prefixed with toast., which controls the behavior of the table's secondary TOAST table, if any (see Section 68.2 for more information about TOAST). If a table parameter value is set and the equivalent toast. parameter is not, the TOAST table will use the table's parameter value. Specifying these parameters for partitioned tables is not supported, but you may specify them for individual leaf partitions.

fillfactor(integer)

The fillfactor for a table is a percentage between 10 and 100. 100 (complete packing) is the default. When a smaller fillfactor is specified, INSERT operations pack table pages only to the indicated percentage; the remaining space on each page is reserved for updating rows on that page. This gives UPDATE a chance to place the updated copy of a row on the same page as the original, which is more efficient than placing it on a different page. For a table whose entries are never updated, complete packing is the best choice, but in heavily updated tables smaller fillfactors are appropriate. This parameter cannot be set for TOAST tables.

toast_tuple_target(integer)

The toast_tuple_target specifies the minimum tuple length required before we try to move long column values into TOAST tables, and is also the target length we try to reduce the length below once toasting begins. This only affects columns marked as either External or Extended and applies only to new tuples - there is no effect on existing rows. By default this parameter is set to allow at least 4 tuples per block, which with the default blocksize will be 2040 bytes. Valid values are between 128 bytes and the (blocksize - header), by default 8160 bytes. Changing this value may not be useful for very short or very long rows. Note that the default setting is often close to optimal, and it is possible that setting this parameter could have negative effects in some cases. This parameter cannot be set for TOAST tables.

parallel_workers(integer)

This sets the number of workers that should be used to assist a parallel scan of this table. If not set, the system will determine a value based on the relation size. The actual number of workers chosen by the planner or by utility statements that use parallel scans may be less, for example due to the setting of max_worker_processes.

autovacuum_enabled, toast.autovacuum_enabled(boolean)

Enables or disables the autovacuum daemon for a particular table. If true, the autovacuum daemon will perform automatic VACUUM and/or ANALYZE operations on this table following the rules discussed in Section 24.1.6. If false, this table will not be autovacuumed, except to prevent transaction ID wraparound. See Section 24.1.5 for more about wraparound prevention. Note that the autovacuum daemon does not run at all (except to prevent transaction ID wraparound) if the autovacuum parameter is false; setting individual tables' storage parameters does not override that. Therefore there is seldom much point in explicitly setting this storage parameter to true, only to false.

autovacuum_vacuum_threshold, toast.autovacuum_vacuum_threshold(integer)

Per-table value for autovacuum_vacuum_threshold parameter.
CREATE TABLE autovacuum_vacuum_scale_factor, toast.autovacuum_vacuum_scale_factor (float4)

Per-table value for autovacuum_vacuum_scale_factor parameter.

autovacuum_analyze_threshold(integer)

Per-table value for autovacuum_analyze_threshold parameter.

autovacuum_analyze_scale_factor(float4)

Per-table value for autovacuum_analyze_scale_factor parameter.

autovacuum_vacuum_cost_delay, toast.autovacuum_vacuum_cost_delay (integer)

Per-table value for autovacuum_vacuum_cost_delay parameter.

autovacuum_vacuum_cost_limit, toast.autovacuum_vacuum_cost_limit (integer)

Per-table value for autovacuum_vacuum_cost_limit parameter.

autovacuum_freeze_min_age, toast.autovacuum_freeze_min_age (integer)

Per-table value for vacuum_freeze_min_age parameter. Note that autovacuum will ignore per-table autovacuum_freeze_min_age parameters that are larger than half the system-wide autovacuum_freeze_max_age setting.

autovacuum_freeze_max_age, toast.autovacuum_freeze_max_age (integer)

Per-table value for autovacuum_freeze_max_age parameter. Note that autovacuum will ignore per-table autovacuum_freeze_max_age parameters that are larger than the system-wide setting (it can only be set smaller).

autovacuum_freeze_table_age, toast.autovacuum_freeze_table_age (integer)

Per-table value for vacuum_freeze_table_age parameter.

autovacuum_multixact_freeze_min_age, toast.autovacuum_multixact_freeze_min_age (integer)

Per-table value for vacuum_multixact_freeze_min_age parameter. Note that autovacuum will ignore per-table autovacuum_multixact_freeze_min_age parameters that are larger than half the system-wide autovacuum_multixact_freeze_max_age setting.

autovacuum_multixact_freeze_max_age, toast.autovacuum_multixact_freeze_max_age (integer)

Per-table value for autovacuum_multixact_freeze_max_age parameter. Note that autovacuum will ignore per-table autovacuum_multixact_freeze_max_age parameters that are larger than the system-wide setting (it can only be set smaller).

autovacuum_multixact_freeze_table_age, toast.autovacuum_multixact_freeze_table_age (integer)

Per-table value for vacuum_multixact_freeze_table_age parameter.

log_autovacuum_min_duration, toast.log_autovacuum_min_duration (integer)

Per-table value for log_autovacuum_min_duration parameter.
CREATE TABLE

user_catalog_table(boolean)

Declare the table as an additional catalog table for purposes of logical replication. See Section 49.6.2 for details. This parameter cannot be set for TOAST tables.

Notes

Using OIDs in new applications is not recommended: where possible, using an identity column or other sequence generator as the table's primary key is preferred. However, if your application does make use of OIDs to identify specific rows of a table, it is recommended to create a unique constraint on the oid column of that table, to ensure that OIDs in the table will indeed uniquely identify rows even after counter wraparound. Avoid assuming that OIDs are unique across tables; if you need a database-wide unique identifier, use the combination of tableoid and row OID for the purpose.

Tip

The use of OIDS=FALSE is not recommended for tables with no primary key, since without either an OID or a unique data key, it is difficult to identify specific rows.

PostgreSQL automatically creates an index for each unique constraint and primary key constraint to enforce uniqueness. Thus, it is not necessary to create an index explicitly for primary key columns. (See CREATE INDEX for more information.)

Unique constraints and primary keys are not inherited in the current implementation. This makes the combination of inheritance and unique constraints rather dysfunctional.

A table cannot have more than 1600 columns. (In practice, the effective limit is usually lower because of tuple-length constraints.)

Examples

Create table films and table distributors:

```
CREATE TABLE films (  
code char(5) CONSTRAINT firstkey PRIMARY KEY,  
title varchar(40) NOT NULL,  
did integer NOT NULL,  
date_prod date,  
kind varchar(10),  
len interval hour to minute  
);  

CREATE TABLE distributors (  
did integer PRIMARY KEY GENERATED BY DEFAULT AS IDENTITY,  
nname varchar(40) NOT NULL CHECK (name <> '')  
);  

CREATE TABLE array_int {  
  vector int[][][]  
};
```

Define a unique table constraint for the table films. Unique table constraints can be defined on one or more columns of the table:
CREATE TABLE films (  
code char(5),
title varchar(40),
did integer,
date_prod date,
kind varchar(10),
len  interval hour to minute,
CONSTRAINT production UNIQUE(date_prod)
);

Define a check column constraint:

CREATE TABLE distributors (  
did integer CHECK (did > 100),
name varchar(40)
);

Define a check table constraint:

CREATE TABLE distributors (  
did integer,
name varchar(40),
CONSTRAINT con1 CHECK (did > 100 AND name <> '')
);

Define a primary key table constraint for the table films:

CREATE TABLE films (  
code char(5),
title varchar(40),
did integer,
date_prod date,
kind varchar(10),
len  interval hour to minute,
CONSTRAINT code_title PRIMARY KEY(code,title)
);

Define a primary key constraint for table distributors. The following two examples are equivalent, the first using the table constraint syntax, the second the column constraint syntax:

CREATE TABLE distributors (  
did integer,
name varchar(40),
PRIMARY KEY(did)
);

CREATE TABLE distributors (  
did integer PRIMARY KEY,
name varchar(40)
);

Assign a literal constant default value for the column name, arrange for the default value of column did to be generated by selecting the next value of a sequence object, and make the default value of modtime be the time at which the row is inserted:
CREATE TABLE distributors (  
  name    varchar(40) DEFAULT 'Luso Films',  
  did     integer DEFAULT nextval('distributors_serial'),  
  modtime timestamp DEFAULT current_timestamp  
);

Define two NOT NULL column constraints on the table distributors, one of which is explicitly given a name:

CREATE TABLE distributors (  
  did     integer CONSTRAINT no_null NOT NULL,  
  name    varchar(40) NOT NULL  
);

Define a unique constraint for the name column:

CREATE TABLE distributors (  
  did     integer,  
  name    varchar(40) UNIQUE  
);

The same, specified as a table constraint:

CREATE TABLE distributors (  
  did     integer,  
  name    varchar(40),  
  UNIQUE(name)  
);

Create the same table, specifying 70% fill factor for both the table and its unique index:

CREATE TABLE distributors (  
  did     integer,  
  name    varchar(40),  
  UNIQUE(name) WITH (fillfactor=70)  
) WITH (fillfactor=70);

Create table circles with an exclusion constraint that prevents any two circles from overlapping:

CREATE TABLE circles (  
  c circle,  
  EXCLUDE USING gist (c WITH &&)  
);

Create table cinemas in tablespace diskvol1:

CREATE TABLE cinemas (  
  id serial,  
  name text,  
  location text  
) TABLESPACE diskvol1;

Create a composite type and a typed table:
CREATE TYPE employee_type AS (name text, salary numeric);

CREATE TABLE employees OF employee_type (
    PRIMARY KEY (name),
    salary WITH OPTIONS DEFAULT 1000
);

Create a range partitioned table:

CREATE TABLE measurement (  
    logdate date not null,  
    peaktemp int,  
    unitsales int  
) PARTITION BY RANGE (logdate);

Create a range partitioned table with multiple columns in the partition key:

CREATE TABLE measurement_year_month (  
    logdate date not null,  
    peaktemp int,  
    unitsales int  
) PARTITION BY RANGE (EXTRACT(YEAR FROM logdate), EXTRACT(MONTH FROM logdate));

Create a list partitioned table:

CREATE TABLE cities (  
    city_id bigserial not null,  
    name text not null,  
    population bigint  
) PARTITION BY LIST (left(lower(name), 1));

Create a hash partitioned table:

CREATE TABLE orders (  
    order_id bigint not null,  
    cust_id bigint not null,  
    status text  
) PARTITION BY HASH (order_id);

Create partition of a range partitioned table:

CREATE TABLE measurement_y2016m07 PARTITION OF measurement (  
    unitsales DEFAULT 0  
) FOR VALUES FROM ('2016-07-01') TO ('2016-08-01');

Create a few partitions of a range partitioned table with multiple columns in the partition key:

CREATE TABLE measurement_ym_older PARTITION OF measurement_year_month  
    FOR VALUES FROM (MINVALUE, MINVALUE) TO (2016, 11);
CREATE TABLE measurement_ym_y2016m11
PARTITION OF measurement_year_month
FOR VALUES FROM (2016, 11) TO (2016, 12);

CREATE TABLE measurement_ym_y2016m12
PARTITION OF measurement_year_month
FOR VALUES FROM (2016, 12) TO (2017, 01);

CREATE TABLE measurement_ym_y2017m01
PARTITION OF measurement_year_month
FOR VALUES FROM (2017, 01) TO (2017, 02);

Create partition of a list partitioned table:

CREATE TABLE cities_ab
PARTITION OF cities (CONSTRAINT city_id_nonzero CHECK (city_id != 0))
FOR VALUES IN ('a', 'b');

Create partition of a list partitioned table that is itself further partitioned and then add a partition to it:

CREATE TABLE cities_ab
PARTITION OF cities (CONSTRAINT city_id_nonzero CHECK (city_id != 0))
FOR VALUES IN ('a', 'b') PARTITION BY RANGE (population);

CREATE TABLE cities_ab_10000_to_100000
PARTITION OF cities_ab FOR VALUES FROM (10000) TO (100000);

Create partitions of a hash partitioned table:

CREATE TABLE orders_p1 PARTITION OF orders
FOR VALUES WITH (MODULUS 4, REMAINDER 0);
CREATE TABLE orders_p2 PARTITION OF orders
FOR VALUES WITH (MODULUS 4, REMAINDER 1);
CREATE TABLE orders_p3 PARTITION OF orders
FOR VALUES WITH (MODULUS 4, REMAINDER 2);
CREATE TABLE orders_p4 PARTITION OF orders
FOR VALUES WITH (MODULUS 4, REMAINDER 3);

Create a default partition:

CREATE TABLE cities_partdef
PARTITION OF cities DEFAULT;

**Compatibility**

The CREATE TABLE command conforms to the SQL standard, with exceptions listed below.

**Temporary Tables**

Although the syntax of CREATE TEMPORARY TABLE resembles that of the SQL standard, the effect is not the same. In the standard, temporary tables are defined just once and automatically exist (starting with empty contents) in every session that needs them. PostgreSQL instead requires each session to
issue its own CREATE TEMPORARY TABLE command for each temporary table to be used. This allows different sessions to use the same temporary table name for different purposes, whereas the standard’s approach constrains all instances of a given temporary table name to have the same table structure.

The standard’s definition of the behavior of temporary tables is widely ignored. PostgreSQL’s behavior on this point is similar to that of several other SQL databases.

The SQL standard also distinguishes between global and local temporary tables, where a local temporary table has a separate set of contents for each SQL module within each session, though its definition is still shared across sessions. Since PostgreSQL does not support SQL modules, this distinction is not relevant in PostgreSQL.

For compatibility’s sake, PostgreSQL will accept the GLOBAL and LOCAL keywords in a temporary table declaration, but they currently have no effect. Use of these keywords is discouraged, since future versions of PostgreSQL might adopt a more standard-compliant interpretation of their meaning.

The ON COMMIT clause for temporary tables also resembles the SQL standard, but has some differences. If the ON COMMIT clause is omitted, SQL specifies that the default behavior is ON COMMIT DELETE ROWS. However, the default behavior in PostgreSQL is ON COMMIT PRESERVE ROWS. The ON COMMIT DROP option does not exist in SQL.

Non-deferred Uniqueness Constraints

When a UNIQUE or PRIMARY KEY constraint is not deferrable, PostgreSQL checks for uniqueness immediately whenever a row is inserted or modified. The SQL standard says that uniqueness should be enforced only at the end of the statement; this makes a difference when, for example, a single command updates multiple key values. To obtain standard-compliant behavior, declare the constraint as DEFERRABLE but not deferred (i.e., INITIALLY IMMEDIATE). Be aware that this can be significantly slower than immediate uniqueness checking.

Column Check Constraints

The SQL standard says that CHECK column constraints can only refer to the column they apply to; only CHECK table constraints can refer to multiple columns. PostgreSQL does not enforce this restriction; it treats column and table check constraints alike.

EXCLUDE Constraint

The EXCLUDE constraint type is a PostgreSQL extension.

NULL “Constraint”

The NULL “constraint” (actually a non-constraint) is a PostgreSQL extension to the SQL standard that is included for compatibility with some other database systems (and for symmetry with the NOT NULL constraint). Since it is the default for any column, its presence is simply noise.

Constraint Naming

The SQL standard says that table and domain constraints must have names that are unique across the schema containing the table or domain. PostgreSQL is laxer: it only requires constraint names to be unique across the constraints attached to a particular table or domain. However, this extra freedom does not exist for index-based constraints (UNIQUE, PRIMARY KEY, and EXCLUDE constraints), because the associated index is named the same as the constraint, and index names must be unique across all relations within the same schema.

Currently, PostgreSQL does not record names for NOT NULL constraints at all, so they are not subject to the uniquenes restriction. This might change in a future release.
**Inheritance**

Multiple inheritance via the `INHERITS` clause is a PostgreSQL language extension. SQL:1999 and later define single inheritance using a different syntax and different semantics. SQL:1999-style inheritance is not yet supported by PostgreSQL.

**Zero-column Tables**

PostgreSQL allows a table of no columns to be created (for example, `CREATE TABLE foo();`). This is an extension from the SQL standard, which does not allow zero-column tables. Zero-column tables are not in themselves very useful, but disallowing them creates odd special cases for `ALTER TABLE DROP COLUMN`, so it seems cleaner to ignore this spec restriction.

**Multiple Identity Columns**

PostgreSQL allows a table to have more than one identity column. The standard specifies that a table can have at most one identity column. This is relaxed mainly to give more flexibility for doing schema changes or migrations. Note that the `INSERT` command supports only one override clause that applies to the entire statement, so having multiple identity columns with different behaviors is not well supported.

**LIKE Clause**

While a `LIKE` clause exists in the SQL standard, many of the options that PostgreSQL accepts for it are not in the standard, and some of the standard's options are not implemented by PostgreSQL.

**WITH Clause**

The `WITH` clause is a PostgreSQL extension; neither storage parameters nor OIDs are in the standard.

**Tablespaces**

The PostgreSQL concept of tablespaces is not part of the standard. Hence, the clauses `TABLESPACE` and `USING INDEX TABLESPACE` are extensions.

**Typed Tables**

Typed tables implement a subset of the SQL standard. According to the standard, a typed table has columns corresponding to the underlying composite type as well as one other column that is the “self-referencing column”. PostgreSQL does not support these self-referencing columns explicitly, but the same effect can be had using the OID feature.

**PARTITION BY Clause**

The `PARTITION BY` clause is a PostgreSQL extension.

**PARTITION OF Clause**

The `PARTITION OF` clause is a PostgreSQL extension.

**See Also**

`ALTER TABLE`, `DROP TABLE`, `CREATE TABLE AS`, `CREATE TABLESPACE`, `CREATE TYPE`
CREATE TABLE AS

CREATE TABLE AS — define a new table from the results of a query

Synopsis

```
CREATE [ [ GLOBAL | LOCAL ] [ TEMPORARY | TEMP ] | UNLOGGED ] TABLE
[ IF NOT EXISTS ] table_name
[ (column_name [, ...] ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ] [ WITH OIDS | WITHOUT OIDS ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]
AS query
[ WITH [ NO ] DATA ]
```

Description

CREATE TABLE AS creates a table and fills it with data computed by a SELECT command. The table columns have the names and data types associated with the output columns of the SELECT (except that you can override the column names by giving an explicit list of new column names).

CREATE TABLE AS bears some resemblance to creating a view, but it is really quite different: it creates a new table and evaluates the query just once to fill the new table initially. The new table will not track subsequent changes to the source tables of the query. In contrast, a view re-evaluates its defining SELECT statement whenever it is queried.

Parameters

GLOBAL or LOCAL

Ignored for compatibility. Use of these keywords is deprecated; refer to CREATE TABLE for details.

TEMPORARY or TEMP

If specified, the table is created as a temporary table. Refer to CREATE TABLE for details.

UNLOGGED

If specified, the table is created as an unlogged table. Refer to CREATE TABLE for details.

IF NOT EXISTS

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Refer to CREATE TABLE for details.

`table_name`

The name (optionally schema-qualified) of the table to be created.

`column_name`

The name of a column in the new table. If column names are not provided, they are taken from the output column names of the query.
CREATE TABLE AS

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for the new table; see Storage Parameters for more information. The WITH clause can also include OIDS=TRUE (or just OIDS) to specify that rows of the new table should have OIDs (object identifiers) assigned to them, or OIDS=FALSE to specify that the rows should not have OIDs. See CREATE TABLE for more information.

WITH OIDS
WITHOUT OIDS

These are obsolescent syntaxes equivalent to WITH (OIDS) and WITH (OIDS=FALSE), respectively. If you wish to give both an OIDS setting and storage parameters, you must use the WITH (...) syntax; see above.

ON COMMIT

The behavior of temporary tables at the end of a transaction block can be controlled using ON COMMIT. The three options are:

PRESERVE ROWS

No special action is taken at the ends of transactions. This is the default behavior.

DELETE ROWS

All rows in the temporary table will be deleted at the end of each transaction block. Essentially, an automatic TRUNCATE is done at each commit.

DROP

The temporary table will be dropped at the end of the current transaction block.

TABLESPACE tablespace_name

The tablespace_name is the name of the tablespace in which the new table is to be created. If not specified, default_tablespace is consulted, or temp_tablespaces if the table is temporary.

query

A SELECT, TABLE, or VALUES command, or an EXECUTE command that runs a prepared SELECT, TABLE, or VALUES query.

WITH [ NO ] DATA

This clause specifies whether or not the data produced by the query should be copied into the new table. If not, only the table structure is copied. The default is to copy the data.

Notes

This command is functionally similar to SELECT INTO, but it is preferred since it is less likely to be confused with other uses of the SELECT INTO syntax. Furthermore, CREATE TABLE AS offers a superset of the functionality offered by SELECT INTO.

The CREATE TABLE AS command allows the user to explicitly specify whether OIDs should be included. If the presence of OIDs is not explicitly specified, the default_with_oids configuration variable is used.

Examples

Create a new table films_recent consisting of only recent entries from the table films:
CREATE TABLE AS

CREATE TABLE films_recent AS
    SELECT * FROM films WHERE date_prod >= '2002-01-01';

To copy a table completely, the short form using the TABLE command can also be used:

CREATE TABLE films2 AS
    TABLE films;

Create a new temporary table films_recent, consisting of only recent entries from the table films, using a prepared statement. The new table has OIDs and will be dropped at commit:

PREPARE recentfilms(date) AS
    SELECT * FROM films WHERE date_prod > $1;
CREATE TEMP TABLE films_recent WITH (OIDS) ON COMMIT DROP AS
    EXECUTE recentfilms('2002-01-01');

Compatibility

CREATE TABLE AS conforms to the SQL standard. The following are nonstandard extensions:

- The standard requires parentheses around the subquery clause; in PostgreSQL, these parentheses are optional.
- In the standard, the WITH [ NO ] DATA clause is required; in PostgreSQL it is optional.
- PostgreSQL handles temporary tables in a way rather different from the standard; see CREATE TABLE for details.
- The WITH clause is a PostgreSQL extension; neither storage parameters nor OIDs are in the standard.
- The PostgreSQL concept of tablespaces is not part of the standard. Hence, the clause TABLESPACE is an extension.

See Also

CREATE MATERIALIZED VIEW, CREATE TABLE, EXECUTE, SELECT, SELECT INTO, VALUES
CREATE TABLESPACE

CREATE TABLESPACE — define a new tablespace

Synopsis

CREATE TABLESPACE _tablespace_name_

[ OWNER { new_owner | CURRENT_USER | SESSION_USER } ]

LOCATION 'directory'

[ WITH ( _tablespace_option_ = value [, ... ] ) ]

Description

CREATE TABLESPACE registers a new cluster-wide tablespace. The tablespace name must be distinct from the name of any existing tablespace in the database cluster.

A tablespace allows superusers to define an alternative location on the file system where the data files containing database objects (such as tables and indexes) can reside.

A user with appropriate privileges can pass _tablespace_name_ to CREATE_DATABASE, CREATE_TABLE, CREATE_INDEX or ADD_CONSTRAINT to have the data files for these objects stored within the specified tablespace.

Warning

A tablespace cannot be used independently of the cluster in which it is defined; see Section 22.6.

Parameters

tablespace_name

The name of a tablespace to be created. The name cannot begin with pg_, as such names are reserved for system tablespaces.

user_name

The name of the user who will own the tablespace. If omitted, defaults to the user executing the command. Only superusers can create tablespaces, but they can assign ownership of tablespaces to non-superusers.

directory

The directory that will be used for the tablespace. The directory should be empty and must be owned by the PostgreSQL system user. The directory must be specified by an absolute path name.

tablespace_option

A tablespace parameter to be set or reset. Currently, the only available parameters are seq_page_cost, random_page_cost and effective_io_concurrency. Setting either value for a particular tablespace will override the planner's usual estimate of the cost of reading pages from tables in that tablespace, as established by the configuration parameters of the same name (see seq_page_cost, random_page_cost, effective_io_concurrency). This may be useful if
CREATE TABLESPACE

One tablespace is located on a disk which is faster or slower than the remainder of the I/O subsystem.

Notes

Tablespaces are only supported on systems that support symbolic links.
CREATE TABLESPACE cannot be executed inside a transaction block.

Examples

Create a tablespace dbspace at /data/dbs:

CREATE TABLESPACE dbspace LOCATION '/data/dbs';

Create a tablespace indexspace at /data/indexes owned by user genevieve:

CREATE TABLESPACE indexspace OWNER genevieve LOCATION '/data/indexes';

Compatibility

CREATE TABLESPACE is a PostgreSQL extension.

See Also

CREATE DATABASE, CREATE TABLE, CREATE INDEX, DROP TABLESPACE, ALTER TABLESPACE
CREATE TEXT SEARCH CONFIGURATION

CREATE TEXT SEARCH CONFIGURATION — define a new text search configuration

Synopsis

CREATE TEXT SEARCH CONFIGURATION name (  
    PARSER = parser_name  |  
    COPY = source_config
  )

Description

CREATE TEXT SEARCH CONFIGURATION creates a new text search configuration. A text search configuration specifies a text search parser that can divide a string into tokens, plus dictionaries that can be used to determine which tokens are of interest for searching.

If only the parser is specified, then the new text search configuration initially has no mappings from token types to dictionaries, and therefore will ignore all words. Subsequent ALTER TEXT SEARCH CONFIGURATION commands must be used to create mappings to make the configuration useful. Alternatively, an existing text search configuration can be copied.

If a schema name is given then the text search configuration is created in the specified schema. Otherwise it is created in the current schema.

The user who defines a text search configuration becomes its owner.

Refer to Chapter 12 for further information.

Parameters

name

The name of the text search configuration to be created. The name can be schema-qualified.

parser_name

The name of the text search parser to use for this configuration.

source_config

The name of an existing text search configuration to copy.

Notes

The PARSER and COPY options are mutually exclusive, because when an existing configuration is copied, its parser selection is copied too.

Compatibility

There is no CREATE TEXT SEARCH CONFIGURATION statement in the SQL standard.

See Also

ALTER TEXT SEARCH CONFIGURATION, DROP TEXT SEARCH CONFIGURATION
CREATE TEXT SEARCH DICTIONARY

CREATE TEXT SEARCH DICTIONARY — define a new text search dictionary

Synopsis

```
CREATE TEXT SEARCH DICTIONARY name (  
    TEMPLATE = template  
    [, option = value [, ... ]]  
  )
```

Description

CREATE TEXT SEARCH DICTIONARY creates a new text search dictionary. A text search dictionary specifies a way of recognizing interesting or uninteresting words for searching. A dictionary depends on a text search template, which specifies the functions that actually perform the work. Typically the dictionary provides some options that control the detailed behavior of the template's functions.

If a schema name is given then the text search dictionary is created in the specified schema. Otherwise it is created in the current schema.

The user who defines a text search dictionary becomes its owner.

Refer to Chapter 12 for further information.

Parameters

- **name**
  
The name of the text search dictionary to be created. The name can be schema-qualified.

- **template**
  
The name of the text search template that will define the basic behavior of this dictionary.

- **option**
  
The name of a template-specific option to be set for this dictionary.

- **value**
  
The value to use for a template-specific option. If the value is not a simple identifier or number, it must be quoted (but you can always quote it, if you wish).

The options can appear in any order.

Examples

The following example command creates a Snowball-based dictionary with a nonstandard list of stop words.

```
CREATE TEXT SEARCH DICTIONARY my_russian (  
    template = snowball,  
    language = russian,  
  
```

```
CREATE TEXT SEARCH DICTIONARY

```sql
stopwords = myrussian
```

**Compatibility**

There is no `CREATE TEXT SEARCH DICTIONARY` statement in the SQL standard.

**See Also**

`ALTER TEXT SEARCH DICTIONARY`, `DROP TEXT SEARCH DICTIONARY`
CREATE TEXT SEARCH PARSER

CREATE TEXT SEARCH PARSER — define a new text search parser

Synopsis

CREATE TEXT SEARCH PARSER name (  
START = start_function ,  
GETTOKEN = gettoken_function ,  
END = end_function ,  
LEXTYPES = lextypes_function  
[, HEADLINE = headline_function ]
)

Description

CREATE TEXT SEARCH PARSER creates a new text search parser. A text search parser defines a method for splitting a text string into tokens and assigning types (categories) to the tokens. A parser is not particularly useful by itself, but must be bound into a text search configuration along with some text search dictionaries to be used for searching.

If a schema name is given then the text search parser is created in the specified schema. Otherwise it is created in the current schema.

You must be a superuser to use CREATE TEXT SEARCH PARSER. (This restriction is made because an erroneous text search parser definition could confuse or even crash the server.)

Refer to Chapter 12 for further information.

Parameters

name

The name of the text search parser to be created. The name can be schema-qualified.

start_function

The name of the start function for the parser.

gettoken_function

The name of the get-next-token function for the parser.

end_function

The name of the end function for the parser.

lextypes_function

The name of the lextypes function for the parser (a function that returns information about the set of token types it produces).

headline_function

The name of the headline function for the parser (a function that summarizes a set of tokens).
The function names can be schema-qualified if necessary. Argument types are not given, since the argument list for each type of function is predetermined. All except the headline function are required.

The arguments can appear in any order, not only the one shown above.

**Compatibility**

There is no `CREATE TEXT SEARCH PARSER` statement in the SQL standard.

**See Also**

`ALTER TEXT SEARCH PARSER`, `DROP TEXT SEARCH PARSER`
CREATE TEXT SEARCH TEMPLATE

CREATE TEXT SEARCH TEMPLATE — define a new text search template

Synopsis

CREATE TEXT SEARCH TEMPLATE name ( 
[ INIT = init_function , ]
LEXIZE = lexize_function
)

Description

CREATE TEXT SEARCH TEMPLATE creates a new text search template. Text search templates define the functions that implement text search dictionaries. A template is not useful by itself, but must be instantiated as a dictionary to be used. The dictionary typically specifies parameters to be given to the template functions.

If a schema name is given then the text search template is created in the specified schema. Otherwise it is created in the current schema.

You must be a superuser to use CREATE TEXT SEARCH TEMPLATE. This restriction is made because an erroneous text search template definition could confuse or even crash the server. The reason for separating templates from dictionaries is that a template encapsulates the “unsafe” aspects of defining a dictionary. The parameters that can be set when defining a dictionary are safe for unprivileged users to set, and so creating a dictionary need not be a privileged operation.

Refer to Chapter 12 for further information.

Parameters

name

The name of the text search template to be created. The name can be schema-qualified.

init_function

The name of the init function for the template.

lexize_function

The name of the lexize function for the template.

The function names can be schema-qualified if necessary. Argument types are not given, since the argument list for each type of function is predetermined. The lexize function is required, but the init function is optional.

The arguments can appear in any order, not only the one shown above.

Compatibility

There is no CREATE TEXT SEARCH TEMPLATE statement in the SQL standard.

See Also

ALTER TEXT SEARCH TEMPLATE, DROP TEXT SEARCH TEMPLATE
CREATE TRANSFORM

CREATE TRANSFORM — define a new transform

Synopsis

CREATE [ OR REPLACE ] TRANSFORM FOR type_name LANGUAGE lang_name (  
FROM SQL WITH FUNCTION from_sql_function_name [ (argument_type  
[, ...]) ]),  
TO SQL WITH FUNCTION to_sql_function_name [ (argument_type  
[, ...]) ]  
);

Description

CREATE TRANSFORM defines a new transform. CREATE OR REPLACE TRANSFORM will either  
create a new transform, or replace an existing definition.

A transform specifies how to adapt a data type to a procedural language. For example, when writing  
a function in PL/Python using the hstore type, PL/Python has no prior knowledge how to present  
hstore values in the Python environment. Language implementations usually default to using the  
text representation, but that is inconvenient when, for example, an associative array or a list would  
be more appropriate.

A transform specifies two functions:

• A “from SQL” function that converts the type from the SQL environment to the language. This  
  function will be invoked on the arguments of a function written in the language.

• A “to SQL” function that converts the type from the language to the SQL environment. This function  
  will be invoked on the return value of a function written in the language.

It is not necessary to provide both of these functions. If one is not specified, the language-specific  
default behavior will be used if necessary. (To prevent a transformation in a certain direction from  
happening at all, you could also write a transform function that always errors out.)

To be able to create a transform, you must own and have USAGE privilege on the type, have USAGE  
privilege on the language, and own and have EXECUTE privilege on the from-SQL and to-SQL func-

Parameters

type_name

The name of the data type of the transform.

lang_name

The name of the language of the transform.

from_sql_function_name[(argument_type [, ...])]  

The name of the function for converting the type from the SQL environment to the language. It  
must take one argument of type internal and return type internal. The actual argument will  
be of the type for the transform, and the function should be coded as if it were. (But it is not allowed
to declare an SQL-level function returning internal without at least one argument of type internal.) The actual return value will be something specific to the language implementation. If no argument list is specified, the function name must be unique in its schema.

\[ \text{to_sql_function_name}[(\text{argument_type} [, ...])] \]

The name of the function for converting the type from the language to the SQL environment. It must take one argument of type internal and return the type that is the type for the transform. The actual argument value will be something specific to the language implementation. If no argument list is specified, the function name must be unique in its schema.

Notes

Use \text{DROP TRANSFORM} to remove transforms.

Examples

To create a transform for type \text{hstore} and language \text{plpythonu}, first set up the type and the language:

\begin{verbatim}
CREATE TYPE hstore ...;
CREATE EXTENSION plpythonu;
\end{verbatim}

Then create the necessary functions:

\begin{verbatim}
CREATE FUNCTION hstore_to_plpython(val internal) RETURNS internal
  LANGUAGE C STRICT IMMUTABLE
  AS ...;

CREATE FUNCTION plpython_to_hstore(val internal) RETURNS hstore
  LANGUAGE C STRICT IMMUTABLE
  AS ...;
\end{verbatim}

And finally create the transform to connect them all together:

\begin{verbatim}
CREATE TRANSFORM FOR hstore LANGUAGE plpythonu (  FROM SQL WITH FUNCTION hstore_to_plpython(internal),
  TO SQL WITH FUNCTION plpython_to_hstore(internal) );
\end{verbatim}

In practice, these commands would be wrapped up in an extension.

The contrib section contains a number of extensions that provide transforms, which can serve as real-world examples.

Compatibility

This form of \text{CREATE TRANSFORM} is a PostgreSQL extension. There is a \text{CREATE TRANSFORM} command in the SQL standard, but it is for adapting data types to client languages. That usage is not supported by PostgreSQL.

See Also

CREATE FUNCTION, CREATE LANGUAGE, CREATE TYPE, DROP TRANSFORM
CREATE TRIGGER

CREATE TRIGGER — define a new trigger

Synopsis

CREATE [ CONSTRAINT ] TRIGGER name { BEFORE | AFTER | INSTEAD OF }
{ event [ OR ... ] }
ON table_name
[ FROM referenced_table_name ]
[ NOT DEFERRABLE | [ DEFERRABLE ] [ INITIALLY IMMEDIATE | INITIAL DEFERRED ] ]
[ REFERENCING { { OLD | NEW } TABLE
[ AS ] transition_relation_name ] [ ... ] ]
[ FOR [ EACH ] { ROW | STATEMENT } ]
[ WHEN ( condition ) ]
EXECUTE { FUNCTION | PROCEDURE } function_name ( arguments )

where event can be one of:

INSERT
UPDATE [ OF column_name [, ... ] ]
DELETE
TRUNCATE

Description

CREATE TRIGGER creates a new trigger. The trigger will be associated with the specified table, view, or foreign table and will execute the specified function function_name when certain operations are performed on that table.

The trigger can be specified to fire before the operation is attempted on a row (before constraints are checked and the INSERT, UPDATE, or DELETE is attempted); or after the operation has completed (after constraints are checked and the INSERT, UPDATE, or DELETE has completed); or instead of the operation (in the case of inserts, updates or deletes on a view). If the trigger fires before or instead of the event, the trigger can skip the operation for the current row, or change the row being inserted (for INSERT and UPDATE operations only). If the trigger fires after the event, all changes, including the effects of other triggers, are “visible” to the trigger.

A trigger that is marked FOR EACH ROW is called once for every row that the operation modifies. For example, a DELETE that affects 10 rows will cause any ON DELETE triggers on the target relation to be called 10 separate times, once for each deleted row. In contrast, a trigger that is marked FOR EACH STATEMENT only executes once for any given operation, regardless of how many rows it modifies (in particular, an operation that modifies zero rows will still result in the execution of any applicable FOR EACH STATEMENT triggers).

Triggers that are specified to fire INSTEAD OF the trigger event must be marked FOR EACH ROW, and can only be defined on views. BEFORE and AFTER triggers on a view must be marked as FOR EACH STATEMENT.

In addition, triggers may be defined to fire for TRUNCATE, though only FOR EACH STATEMENT.

The following table summarizes which types of triggers may be used on tables, views, and foreign tables:
CREATE TRIGGER

When | Event       | Row-level                          | Statement-level                        |
-----|-------------|------------------------------------|----------------------------------------|
     | BEFORE      | INSERT/UPDATE/DELETE               | Tables and foreign tables              |
     |             | TRUNCATE                           | —                                      |
     | AFTER       | INSERT/UPDATE/DELETE               | Tables and foreign tables              |
     |             | TRUNCATE                           | —                                      |
     | INSTEAD OF  | INSERT/UPDATE/DELETE               | Views                                  |
     |             | TRUNCATE                           | —                                      |

Also, a trigger definition can specify a Boolean WHEN condition, which will be tested to see whether the trigger should be fired. In row-level triggers the WHEN condition can examine the old and/or new values of columns of the row. Statement-level triggers can also have WHEN conditions, although the feature is not so useful for them since the condition cannot refer to any values in the table.

If multiple triggers of the same kind are defined for the same event, they will be fired in alphabetical order by name.

When the CONSTRAINT option is specified, this command creates a constraint trigger. This is the same as a regular trigger except that the timing of the trigger firing can be adjusted using SET CONSTRAINTS. Constraint triggers must be AFTER ROW triggers on plain tables (not foreign tables). They can be fired either at the end of the statement causing the triggering event, or at the end of the containing transaction; in the latter case they are said to be deferred. A pending deferred-trigger firing can also be forced to happen immediately by using SET CONSTRAINTS. Constraint triggers are expected to raise an exception when the constraints they implement are violated.

The REFERENCING option enables collection of transition relations, which are row sets that include all of the rows inserted, deleted, or modified by the current SQL statement. This feature lets the trigger see a global view of what the statement did, not just one row at a time. This option is only allowed for an AFTER trigger that is not a constraint trigger; also, if the trigger is an UPDATE trigger, it must not specify a column_name list. OLD TABLE may only be specified once, and only for a trigger that can fire on UPDATE or DELETE; it creates a transition relation containing the before-images of all rows updated or deleted by the statement. Similarly, NEW TABLE may only be specified once, and only for a trigger that can fire on UPDATE or INSERT; it creates a transition relation containing the after-images of all rows updated or inserted by the statement.

SELECT does not modify any rows so you cannot create SELECT triggers. Rules and views may provide workable solutions to problems that seem to need SELECT triggers.

Refer to Chapter 39 for more information about triggers.

Parameters

name

The name to give the new trigger. This must be distinct from the name of any other trigger for the same table. The name cannot be schema-qualified — the trigger inherits the schema of its table. For a constraint trigger, this is also the name to use when modifying the trigger's behavior using SET CONSTRAINTS.

BEFORE
AFTER
INSTEAD OF

Determines whether the function is called before, after, or instead of the event. A constraint trigger can only be specified as AFTER.
One of INSERT, UPDATE, DELETE, or TRUNCATE; this specifies the event that will fire the trigger. Multiple events can be specified using OR, except when transition relations are requested.

For UPDATE events, it is possible to specify a list of columns using this syntax:

```
UPDATE OF column_name1 [, column_name2 ... ]
```

The trigger will only fire if at least one of the listed columns is mentioned as a target of the UPDATE command.

INSTEAD OF UPDATE events do not allow a list of columns. A column list cannot be specified when requesting transition relations, either.

table_name

The name (optionally schema-qualified) of the table, view, or foreign table the trigger is for.

referenced_table_name

The (possibly schema-qualified) name of another table referenced by the constraint. This option is used for foreign-key constraints and is not recommended for general use. This can only be specified for constraint triggers.

DEFERRABLE
NOT DEFERRABLE
INITIALLY IMMEDIATE
INITIALLY DEFERRED

The default timing of the trigger. See the CREATE TABLE documentation for details of these constraint options. This can only be specified for constraint triggers.

REFERENCING

This keyword immediately precedes the declaration of one or two relation names that provide access to the transition relations of the triggering statement.

OLD TABLE
NEW TABLE

This clause indicates whether the following relation name is for the before-image transition relation or the after-image transition relation.

transition_relation_name

The (unqualified) name to be used within the trigger for this transition relation.

FOR EACH ROW
FOR EACH STATEMENT

This specifies whether the trigger function should be fired once for every row affected by the trigger event, or just once per SQL statement. If neither is specified, FOR EACH STATEMENT is the default. Constraint triggers can only be specified FOR EACH ROW.

condition

A Boolean expression that determines whether the trigger function will actually be executed. If WHEN is specified, the function will only be called if the condition returns true. In FOR EACH ROW triggers, the WHEN condition can refer to columns of the old and/or new row val-
CREATE TRIGGER

uses by writing OLD.column_name or NEW.column_name respectively. Of course, INSERT triggers cannot refer to OLD and DELETE triggers cannot refer to NEW.

INSTEAD_OF triggers do not support WHEN conditions.

Currently, WHEN expressions cannot contain subqueries.

Note that for constraint triggers, evaluation of the WHEN condition is not deferred, but occurs immediately after the row update operation is performed. If the condition does not evaluate to true then the trigger is not queued for deferred execution.

function_name

A user-supplied function that is declared as taking no arguments and returning type trigger, which is executed when the trigger fires.

In the syntax of CREATE TRIGGER, the keywords FUNCTION and PROCEDURE are equivalent, but the referenced function must in any case be a function, not a procedure. The use of the keyword PROCEDURE here is historical and deprecated.

arguments

An optional comma-separated list of arguments to be provided to the function when the trigger is executed. The arguments are literal string constants. Simple names and numeric constants can be written here, too, but they will all be converted to strings. Please check the description of the implementation language of the trigger function to find out how these arguments can be accessed within the function; it might be different from normal function arguments.

Notes

To create a trigger on a table, the user must have the TRIGGER privilege on the table. The user must also have EXECUTE privilege on the trigger function.

Use DROP TRIGGER to remove a trigger.

A column-specific trigger (one defined using the UPDATE OF column_name syntax) will fire when any of its columns are listed as targets in the UPDATE command's SET list. It is possible for a column's value to change even when the trigger is not fired, because changes made to the row's contents by BEFORE UPDATE triggers are not considered. Conversely, a command such as UPDATE ... SET x = x ... will fire a trigger on column x, even though the column's value did not change.

In a BEFORE trigger, the WHEN condition is evaluated just before the function is or would be executed, so using WHEN is not materially different from testing the same condition at the beginning of the trigger function. Note in particular that the NEW row seen by the condition is the current value, as possibly modified by earlier triggers. Also, a BEFORE trigger's WHEN condition is not allowed to examine the system columns of the NEW row (such as oid), because those won't have been set yet.

In an AFTER trigger, the WHEN condition is evaluated just after the row update occurs, and it determines whether an event is queued to fire the trigger at the end of statement. So when an AFTER trigger's WHEN condition does not return true, it is not necessary to queue an event nor to re-fetch the row at end of statement. This can result in significant speedups in statements that modify many rows, if the trigger only needs to be fired for a few of the rows.

In some cases it is possible for a single SQL command to fire more than one kind of trigger. For instance an INSERT with an ON CONFLICT DO UPDATE clause may cause both insert and update operations, so it will fire both kinds of triggers as needed. The transition relations supplied to triggers are specific to their event type; thus an INSERT trigger will see only the inserted rows, while an UPDATE trigger will see only the updated rows.

Row updates or deletions caused by foreign-key enforcement actions, such as ON UPDATE CASCADE or ON DELETE SET NULL, are treated as part of the SQL command that caused them (note that such
actions are never deferred). Relevant triggers on the affected table will be fired, so that this provides another way in which a SQL command might fire triggers not directly matching its type. In simple cases, triggers that request transition relations will see all changes caused in their table by a single original SQL command as a single transition relation. However, there are cases in which the presence of an AFTER ROW trigger that requests transition relations will cause the foreign-key enforcement actions triggered by a single SQL command to be split into multiple steps, each with its own transition relation(s). In such cases, any statement-level triggers that are present will be fired once per creation of a transition relation set, ensuring that the triggers see each affected row in a transition relation once and only once.

Statement-level triggers on a view are fired only if the action on the view is handled by a row-level INSTEAD OF trigger. If the action is handled by an INSTEAD rule, then whatever statements are emitted by the rule are executed in place of the original statement naming the view, so that the triggers that will be fired are those on tables named in the replacement statements. Similarly, if the view is automatically updatable, then the action is handled by automatically rewriting the statement into an action on the view's base table, so that the base table's statement-level triggers are the ones that are fired.

Creating a row-level trigger on a partitioned table will cause identical triggers to be created in all its existing partitions; and any partitions created or attached later will contain an identical trigger, too. Triggers on partitioned tables may only be AFTER.

Modifying a partitioned table or a table with inheritance children fires statement-level triggers attached to the explicitly named table, but not statement-level triggers for its partitions or child tables. In contrast, row-level triggers are fired on the rows in affected partitions or child tables, even if they are not explicitly named in the query. If a statement-level trigger has been defined with transition relations named by a REFERENCING clause, then before and after images of rows are visible from all affected partitions or child tables. In the case of inheritance children, the row images include only columns that are present in the table that the trigger is attached to. Currently, row-level triggers with transition relations cannot be defined on partitions or inheritance child tables.

In PostgreSQL versions before 7.3, it was necessary to declare trigger functions as returning the placeholder type opaque, rather than trigger. To support loading of old dump files, CREATE TRIGGER will accept a function declared as returning opaque, but it will issue a notice and change the function's declared return type to trigger.

### Examples

Execute the function check_account_update whenever a row of the table accounts is about to be updated:

```sql
CREATE TRIGGER check_update
BEFORE UPDATE ON accounts
FOR EACH ROW
EXECUTE FUNCTION check_account_update();
```

The same, but only execute the function if column balance is specified as a target in the UPDATE command:

```sql
CREATE TRIGGER check_update
BEFORE UPDATE OF balance ON accounts
FOR EACH ROW
EXECUTE FUNCTION check_account_update();
```

This form only executes the function if column balance has in fact changed value:
CREATE TRIGGER check_update
BEFORE UPDATE ON accounts
FOR EACH ROW
WHEN (OLD.balance IS DISTINCT FROM NEW.balance)
EXECUTE FUNCTION check_account_update();

Call a function to log updates of accounts, but only if something changed:

CREATE TRIGGER log_update
AFTER UPDATE ON accounts
FOR EACH ROW
WHEN (OLD.* IS DISTINCT FROM NEW.*)
EXECUTE FUNCTION log_account_update();

Execute the function view_insert_row for each row to insert rows into the tables underlying a view:

CREATE TRIGGER view_insert
    INSTEAD OF INSERT ON my_view
    FOR EACH ROW
    EXECUTE FUNCTION view_insert_row();

Execute the function check_transfer_balances_to_zero for each statement to confirm that the transfer rows offset to a net of zero:

CREATE TRIGGER transfer_insert
AFTER INSERT ON transfer
REFERENCING NEW TABLE AS inserted
FOR EACH STATEMENT
EXECUTE FUNCTION check_transfer_balances_to_zero();

Execute the function check_matching_pairs for each row to confirm that changes are made to matching pairs at the same time (by the same statement):

CREATE TRIGGER paired_items_update
AFTER UPDATE ON paired_items
REFERENCING NEW TABLE AS newtab OLD TABLE AS oldtab
FOR EACH ROW
EXECUTE FUNCTION check_matching_pairs();

Section 39.4 contains a complete example of a trigger function written in C.

Compatibility

The CREATE TRIGGER statement in PostgreSQL implements a subset of the SQL standard. The following functionalities are currently missing:

- While transition table names for AFTER triggers are specified using the REFERENCING clause in the standard way, the row variables used in FOR EACH ROW triggers may not be specified in a REFERENCING clause. They are available in a manner that is dependent on the language in which the trigger function is written, but is fixed for any one language. Some languages effectively behave as though there is a REFERENCING clause containing OLD ROW AS OLD NEW ROW AS NEW.

- The standard allows transition tables to be used with column-specific UPDATE triggers, but then the set of rows that should be visible in the transition tables depends on the trigger's column list. This is not currently implemented by PostgreSQL.
CREATE TRIGGER

- PostgreSQL only allows the execution of a user-defined function for the triggered action. The standard allows the execution of a number of other SQL commands, such as CREATE TABLE, as the triggered action. This limitation is not hard to work around by creating a user-defined function that executes the desired commands.

SQL specifies that multiple triggers should be fired in time-of-creation order. PostgreSQL uses name order, which was judged to be more convenient.

SQL specifies that BEFORE DELETE triggers on cascaded deletes fire after the cascaded DELETE completes. The PostgreSQL behavior is for BEFORE DELETE to always fire before the delete action, even a cascading one. This is considered more consistent. There is also nonstandard behavior if BEFORE triggers modify rows or prevent updates during an update that is caused by a referential action. This can lead to constraint violations or stored data that does not honor the referential constraint.

The ability to specify multiple actions for a single trigger using OR is a PostgreSQL extension of the SQL standard.

The ability to fire triggers for TRUNCATE is a PostgreSQL extension of the SQL standard, as is the ability to define statement-level triggers on views.

CREATE CONSTRAINT TRIGGER is a PostgreSQL extension of the SQL standard.

**See Also**

ALTER TRIGGER, DROP TRIGGER, CREATE FUNCTION, SET CONSTRAINTS
CREATE TYPE

CREATE TYPE — define a new data type

Synopsis

CREATE TYPE name AS
( [ attribute_name data_type [ COLLATE collation ] [, ... ] ] )

CREATE TYPE name AS ENUM
( [ 'label' [, ... ] ] )

CREATE TYPE name AS RANGE (
  SUBTYPE = subtype
  [, SUBTYPE_OPCLASS = subtype_operator_class ]
  [, COLLATION = collation ]
  [, CANONICAL = canonical_function ]
  [, SUBTYPE_DIFF = subtype_diff_function ]
)

CREATE TYPE name (
  INPUT = input_function,
  OUTPUT = output_function
  [, RECEIVE = receive_function ]
  [, SEND = send_function ]
  [, TYPMOD_IN = type_modifier_input_function ]
  [, TYPMOD_OUT = type_modifier_output_function ]
  [, ANALYZE = analyze_function ]
  [, INTERNALLENGTH = { internallength | VARIABLE } ]
  [, PASSEDBYVALUE ]
  [, ALIGNMENT = alignment ]
  [, STORAGE = storage ]
  [, LIKE = like_type ]
  [, CATEGORY = category ]
  [, PREFERRED = preferred ]
  [, DEFAULT = default ]
  [, ELEMENT = element ]
  [, DELIMITER = delimiter ]
  [, COLLATABLE = collatable ]
)

Description

CREATE TYPE registers a new data type for use in the current database. The user who defines a type becomes its owner.

If a schema name is given then the type is created in the specified schema. Otherwise it is created in the current schema. The type name must be distinct from the name of any existing type or domain in the same schema. (Because tables have associated data types, the type name must also be distinct from the name of any existing table in the same schema.)

There are five forms of CREATE TYPE, as shown in the syntax synopsis above. They respectively create a composite type, an enum type, a range type, a base type, or a shell type. The first four of
these are discussed in turn below. A shell type is simply a placeholder for a type to be defined later; it is created by issuing `CREATE TYPE` with no parameters except for the type name. Shell types are needed as forward references when creating range types and base types, as discussed in those sections.

**Composite Types**

The first form of `CREATE TYPE` creates a composite type. The composite type is specified by a list of attribute names and data types. An attribute's collation can be specified too, if its data type is collatable. A composite type is essentially the same as the row type of a table, but using `CREATE TYPE` avoids the need to create an actual table when all that is wanted is to define a type. A standalone composite type is useful, for example, as the argument or return type of a function.

To be able to create a composite type, you must have `USAGE` privilege on all attribute types.

**Enumerated Types**

The second form of `CREATE TYPE` creates an enumerated (enum) type, as described in Section 8.7. Enum types take a list of one or more quoted labels, each of which must be less than `NAMEDATALEN` bytes long (64 bytes in a standard PostgreSQL build).

**Range Types**

The third form of `CREATE TYPE` creates a new range type, as described in Section 8.17. The range type's `subtype` can be any type with an associated b-tree operator class (to determine the ordering of values for the range type). Normally the subtype's default b-tree operator class is used to determine ordering; to use a non-default operator class, specify its name with `subtype_opclass`. If the subtype is collatable, and you want to use a non-default collation in the range's ordering, specify the desired collation with the `collation` option.

The optional `canonical` function must take one argument of the range type being defined, and return a value of the same type. This is used to convert range values to a canonical form, when applicable. See Section 8.17.8 for more information. Creating a canonical function is a bit tricky, since it must be defined before the range type can be declared. To do this, you must first create a shell type, which is a placeholder type that has no properties except a name and an owner. This is done by issuing the command `CREATE TYPE name`, with no additional parameters. Then the function can be declared using the shell type as argument and result, and finally the range type can be declared using the same name. This automatically replaces the shell type entry with a valid range type.

The optional `subtype_diff` function must take two values of the `subtype` type as argument, and return a `double precision` value representing the difference between the two given values. While this is optional, providing it allows much greater efficiency of GiST indexes on columns of the range type. See Section 8.17.8 for more information.

**Base Types**

The fourth form of `CREATE TYPE` creates a new base type (scalar type). To create a new base type, you must be a superuser. (This restriction is made because an erroneous type definition could confuse or even crash the server.)

The parameters can appear in any order, not only that illustrated above, and most are optional. You must register two or more functions (using `CREATE FUNCTION`) before defining the type. The support functions `input_function` and `output_function` are required, while the functions `receive_function`, `send_function`, `type_modifier_input_function`, `type_modifier_output_function` and `analyze_function` are optional. Generally these functions have to be coded in C or another low-level language.

The `input_function` converts the type's external textual representation to the internal representation used by the operators and functions defined for the type. `output_function` performs the reverse transformation. The input function can be declared as taking one argument of type `cstring`, or as taking three arguments of types `cstring`, `oid`, `integer`. The first argument is the input text
as a C string, the second argument is the type’s own OID (except for array types, which instead receive their element type’s OID), and the third is the typmod of the destination column, if known (-1 will be passed if not). The input function must return a value of the data type itself. Usually, an input function should be declared STRICT; if it is not, it will be called with a NULL first parameter when reading a NULL input value. The function must still return NULL in this case, unless it raises an error. (This case is mainly meant to support domain input functions, which might need to reject NULL inputs.) The output function must be declared as taking one argument of the new data type. The output function must return type cstring. Output functions are not invoked for NULL values.

The optional receive_function converts the type’s external binary representation to the internal representation. If this function is not supplied, the type cannot participate in binary input. The binary representation should be chosen to be cheap to convert to internal form, while being reasonably portable. (For example, the standard integer data types use network byte order as the external binary representation, while the internal representation is in the machine’s native byte order.) The receive function should perform adequate checking to ensure that the value is valid. The receive function can be declared as taking one argument of type internal, or as taking three arguments of types internal, oid, integer. The first argument is a pointer to a StringInfo buffer holding the received byte string; the optional arguments are the same as for the text input function. The receive function must return a value of the data type itself. Usually, a receive function should be declared STRICT; if it is not, it will be called with a NULL first parameter when reading a NULL input value. The function must still return NULL in this case, unless it raises an error. (This case is mainly meant to support domain receive functions, which might need to reject NULL inputs.) Similarly, the optional send_function converts from the internal representation to the external binary representation. If this function is not supplied, the type cannot participate in binary output. The send function must be declared as taking one argument of the new data type. The send function must return type bytea. Send functions are not invoked for NULL values.

You should at this point be wondering how the input and output functions can be declared to have results or arguments of the new type, when they have to be created before the new type can be created. The answer is that the type should first be defined as a shell type, which is a placeholder type that has no properties except a name and an owner. This is done by issuing the command CREATE TYPE name, with no additional parameters. Then the C I/O functions can be defined referencing the shell type. Finally, CREATE TYPE with a full definition replaces the shell entry with a complete, valid type definition, after which the new type can be used normally.

The optional type_modifier_input_function and type_modifier_output_function are needed if the type supports modifiers, that is optional constraints attached to a type declaration, such as char(5) or numeric(30,2). PostgreSQL allows user-defined types to take one or more simple constants or identifiers as modifiers. However, this information must be capable of being packed into a single non-negative integer value for storage in the system catalogs. The type modifier_input_function is passed the declared modifier(s) in the form of a cstring array. It must check the values for validity (throwing an error if they are wrong), and if they are correct, return a single non-negative integer value that will be stored as the column “typmod”. Type modifiers will be rejected if the type does not have a type Modifier input function. The type modifier_output_function converts the internal integer typmod value back to the correct form for user display. It must return a cstring value that is the exact string to append to the type name; for example numeric’s function might return (30,2). It is allowed to omit the type Modifier output_function, in which case the default display format is just the stored typmod integer value enclosed in parentheses.

The optional analyze_function performs type-specific statistics collection for columns of the data type. By default, ANALYZE will attempt to gather statistics using the type’s “equals” and “less-than” operators, if there is a default b-tree operator class for the type. For non-scalar types this behavior is likely to be unsuitable, so it can be overridden by specifying a custom analysis function. The analysis function must be declared to take a single argument of type internal, and return a boolean result. The detailed API for analysis functions appears in src/include/commands/vacuum.h.

While the details of the new type’s internal representation are only known to the I/O functions and other functions you create to work with the type, there are several properties of the internal representation
that must be declared to PostgreSQL. Foremost of these is \textit{internallength}. Base data types can be fixed-length, in which case \textit{internallength} is a positive integer, or variable-length, indicated by setting \textit{internallength} to \texttt{VARIABLE}. (Internally, this is represented by setting \texttt{typlen} to -1.) The internal representation of all variable-length types must start with a 4-byte integer giving the total length of this value of the type. (Note that the length field is often encoded, as described in Section 68.2; it's unwise to access it directly.)

The optional flag \texttt{PASSEDBYVALUE} indicates that values of this data type are passed by value, rather than by reference. Types passed by value must be fixed-length, and their internal representation cannot be larger than the size of the \texttt{Datum} type (4 bytes on some machines, 8 bytes on others).

The \textit{alignment} parameter specifies the storage alignment required for the data type. The allowed values equate to alignment on 1, 2, 4, or 8 byte boundaries. Note that variable-length types must have an alignment of at least 4, since they necessarily contain an \texttt{int4} as their first component.

The \textit{storage} parameter allows selection of storage strategies for variable-length data types. (Only \texttt{plain} is allowed for fixed-length types.) \texttt{plain} specifies that data of the type will always be stored in-line and not compressed. \texttt{extended} specifies that the system will first try to compress a long data value, and will move the value out of the main table row if it's still too long. \texttt{external} allows the value to be moved out of the main table, but the system will not try to compress it. \texttt{main} allows compression, but discourages moving the value out of the main table. (Data items with this storage strategy might still be moved out of the main table if there is no other way to make a row fit, but they will be kept in the main table preferentially over \texttt{extended} and \texttt{external} items.)

All \textit{storage} values other than \texttt{plain} imply that the functions of the data type can handle values that have been \textit{toasted}, as described in Section 68.2 and Section 38.12.1. The specific other value given merely determines the default TOAST storage strategy for columns of a toastable data type; users can pick other strategies for individual columns using \texttt{ALTER TABLE SET STORAGE}.

The \textit{like_type} parameter provides an alternative method for specifying the basic representation properties of a data type: copy them from some existing type. The values of \textit{internallength}, \texttt{PASSEDBYVALUE}, \textit{alignment}, and \textit{storage} are copied from the named type. (It is possible, though usually undesirable, to override some of these values by specifying them along with the \texttt{LIKE} clause.) Specifying representation this way is especially useful when the low-level implementation of the new type “piggybacks” on an existing type in some fashion.

The \textit{category} and \textit{preferred} parameters can be used to help control which implicit cast will be applied in ambiguous situations. Each data type belongs to a category named by a single ASCII character, and each type is either “preferred” or not within its category. The parser will prefer casting to preferred types (but only from other types within the same category) when this rule is helpful in resolving overloaded functions or operators. For more details see Chapter 10. For types that have no implicit casts to or from any other types, it is sufficient to leave these settings at the defaults. However, for a group of related types that have implicit casts, it is often helpful to mark them all as belonging to a category and select one or two of the “most general” types as being preferred within the category. The \textit{category} parameter is especially useful when adding a user-defined type to an existing built-in category, such as the numeric or string types. However, it is also possible to create new entirely-user-defined type categories. Select any ASCII character other than an upper-case letter to name such a category.

A default value can be specified, in case a user wants columns of the data type to default to something other than the null value. Specify the default with the \texttt{DEFAULT} key word. (Such a default can be overridden by an explicit \texttt{DEFAULT} clause attached to a particular column.)

To indicate that a type is an array, specify the type of the array elements using the \texttt{ELEMENT} key word. For example, to define an array of 4-byte integers (\texttt{int4}), specify \texttt{ELEMENT = int4}. More details about array types appear below.

To indicate the delimiter to be used between values in the external representation of arrays of this type, \textit{delimiter} can be set to a specific character. The default delimiter is the comma (,). Note that the delimiter is associated with the array element type, not the array type itself.
CREATE TYPE

If the optional Boolean parameter collatable is true, column definitions and expressions of the type may carry collation information through use of the COLLATE clause. It is up to the implementations of the functions operating on the type to actually make use of the collation information; this does not happen automatically merely by marking the type collatable.

Array Types

Whenever a user-defined type is created, PostgreSQL automatically creates an associated array type, whose name consists of the element type's name prepended with an underscore, and truncated if necessary to keep it less than NAMEDATALEN bytes long. (If the name so generated collides with an existing type name, the process is repeated until a non-colliding name is found.) This implicitly-created array type is variable length and uses the built-in input and output functions array_in and array_out. The array type tracks any changes in its element type's owner or schema, and is dropped if the element type is.

You might reasonably ask why there is an ELEMENT option, if the system makes the correct array type automatically. The only case where it's useful to use ELEMENT is when you are making a fixed-length type that happens to be internally an array of a number of identical things, and you want to allow these things to be accessed directly by subscripting, in addition to whatever operations you plan to provide for the type as a whole. For example, type point is represented as just two floating-point numbers, which can be accessed using point[0] and point[1]. Note that this facility only works for fixed-length types whose internal form is exactly a sequence of identical fixed-length fields. A subscriptable variable-length type must have the generalized internal representation used by array_in and array_out. For historical reasons (i.e., this is clearly wrong but it's far too late to change it), subscripting of fixed-length array types starts from zero, rather than from one as for variable-length arrays.

Parameters

name

The name (optionally schema-qualified) of a type to be created.

attribute_name

The name of an attribute (column) for the composite type.

data_type

The name of an existing data type to become a column of the composite type.

collation

The name of an existing collation to be associated with a column of a composite type, or with a range type.

label

A string literal representing the textual label associated with one value of an enum type.

subtype

The name of the element type that the range type will represent ranges of.

subtype_operator_class

The name of a b-tree operator class for the subtype.

canonical_function

The name of the canonicalization function for the range type.
CREATE TYPE

subtype_diff_function
The name of a difference function for the subtype.

input_function
The name of a function that converts data from the type's external textual form to its internal form.

output_function
The name of a function that converts data from the type's internal form to its external textual form.

receive_function
The name of a function that converts data from the type's external binary form to its internal form.

send_function
The name of a function that converts data from the type's internal form to its external binary form.

type_modifier_input_function
The name of a function that converts an array of modifier(s) for the type into internal form.

type_modifier_output_function
The name of a function that converts the internal form of the type's modifier(s) to external textual form.

analyze_function
The name of a function that performs statistical analysis for the data type.

internallength
A numeric constant that specifies the length in bytes of the new type's internal representation. The default assumption is that it is variable-length.

alignment
The storage alignment requirement of the data type. If specified, it must be char, int2, int4, or double; the default is int4.

storage
The storage strategy for the data type. If specified, must be plain, external, extended, or main; the default is plain.

like_type
The name of an existing data type that the new type will have the same representation as. The values of internallength, passedbyvalue, alignment, and storage are copied from that type, unless overridden by explicit specification elsewhere in this CREATE TYPE command.

category
The category code (a single ASCII character) for this type. The default is 'U' for "user-defined type". Other standard category codes can be found in Table 52.63. You may also choose other ASCII characters in order to create custom categories.

preferred
True if this type is a preferred type within its type category, else false. The default is false. Be very careful about creating a new preferred type within an existing type category, as this could cause surprising changes in behavior.
CREATE TYPE

default

The default value for the data type. If this is omitted, the default is null.

element

The type being created is an array; this specifies the type of the array elements.

delimiter

The delimiter character to be used between values in arrays made of this type.

collatable

True if this type's operations can use collation information. The default is false.

Notes

Because there are no restrictions on use of a data type once it's been created, creating a base type or range type is tantamount to granting public execute permission on the functions mentioned in the type definition. This is usually not an issue for the sorts of functions that are useful in a type definition. But you might want to think twice before designing a type in a way that would require “secret” information to be used while converting it to or from external form.

Before PostgreSQL version 8.3, the name of a generated array type was always exactly the element type's name with one underscore character (_) prepended. (Type names were therefore restricted in length to one less character than other names.) While this is still usually the case, the array type name may vary from this in case of maximum-length names or collisions with user type names that begin with underscore. Writing code that depends on this convention is therefore deprecated. Instead, use pg_type.typarray to locate the array type associated with a given type.

It may be advisable to avoid using type and table names that begin with underscore. While the server will change generated array type names to avoid collisions with user-given names, there is still risk of confusion, particularly with old client software that may assume that type names beginning with underscores always represent arrays.

Before PostgreSQL version 8.2, the shell-type creation syntax CREATE TYPE name did not exist. The way to create a new base type was to create its input function first. In this approach, PostgreSQL will first see the name of the new data type as the return type of the input function. The shell type is implicitly created in this situation, and then it can be referenced in the definitions of the remaining I/O functions. This approach still works, but is deprecated and might be disallowed in some future release. Also, to avoid accidentally cluttering the catalogs with shell types as a result of simple typos in function definitions, a shell type will only be made this way when the input function is written in C.

In PostgreSQL versions before 7.3, it was customary to avoid creating a shell type at all, by replacing the functions' forward references to the type name with the placeholder pseudo-type opaque. The cstring arguments and results also had to be declared as opaque before 7.3. To support loading of old dump files, CREATE TYPE will accept I/O functions declared using opaque, but it will issue a notice and change the function declarations to use the correct types.

Examples

This example creates a composite type and uses it in a function definition:

CREATE TYPE compfoo AS (f1 int, f2 text);

CREATE FUNCTION getfoo() RETURNS SETOF compfoo AS $$
    SELECT fooid, fooname FROM foo
$$
CREATE TYPE $$ LANGUAGE SQL;

This example creates an enumerated type and uses it in a table definition:

CREATE TYPE bug_status AS ENUM ('new', 'open', 'closed');

CREATE TABLE bug (  
id serial,  
description text,  
status bug_status
);

This example creates a range type:

CREATE TYPE float8_range AS RANGE (subtype = float8, subtype_diff = float8mi);

This example creates the base data type box and then uses the type in a table definition:

CREATE TYPE box;

CREATE FUNCTION my_box_in_function(cstring) RETURNS box AS ... ;
CREATE FUNCTION my_box_out_function(box) RETURNS cstring AS ... ;

CREATE TYPE box (  
  INTERNALLENGTH = 16,  
  INPUT = my_box_in_function,  
  OUTPUT = my_box_out_function
);

CREATE TABLE myboxes (  
id integer,  
description box
);

If the internal structure of box were an array of four float4 elements, we might instead use:

CREATE TYPE box (  
  INTERNALLENGTH = 16,  
  INPUT = my_box_in_function,  
  OUTPUT = my_box_out_function,  
  ELEMENT = float4
);

which would allow a box value's component numbers to be accessed by subscripting. Otherwise the type behaves the same as before.

This example creates a large object type and uses it in a table definition:

CREATE TYPE bigobj (  
  INPUT = lo_filein, OUTPUT = lo_fileout,  
  INTERNALLENGTH = VARIABLE
);
CREATE TABLE big_objs (  
id integer,
...
CREATE TYPE obj bigobj

More examples, including suitable input and output functions, are in Section 38.12.

Compatibility

The first form of the CREATE TYPE command, which creates a composite type, conforms to the SQL standard. The other forms are PostgreSQL extensions. The CREATE TYPE statement in the SQL standard also defines other forms that are not implemented in PostgreSQL.

The ability to create a composite type with zero attributes is a PostgreSQL-specific deviation from the standard (analogous to the same case in CREATE TABLE).

See Also

ALTER TYPE, CREATE DOMAIN, CREATE FUNCTION, DROP TYPE
CREATE USER

CREATE USER — define a new database role

Synopsis

CREATE USER name [ [ WITH ] option [ ... ] ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICATION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT connlimit
- [ ENCRYPTED ] PASSWORD 'password'
- VALID UNTIL 'timestamp'
- IN ROLE role_name [, ...]
- IN GROUP role_name [, ...]
- ROLE role_name [, ...]
- ADMIN role_name [, ...]
- USER role_name [, ...]
- SYSID uid

Description

CREATE USER is now an alias for CREATE ROLE. The only difference is that when the command is spelled CREATE USER, LOGIN is assumed by default, whereas NOLOGIN is assumed when the command is spelled CREATE ROLE.

Compatibility

The CREATE USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also

CREATE ROLE
CREATE USER MAPPING

CREATE USER MAPPING — define a new mapping of a user to a foreign server

Synopsis

CREATE USER MAPPING [ IF NOT EXISTS ] FOR { user_name | CURRENT_USER | PUBLIC } SERVER server_name
[ OPTIONS ( option 'value' [ , ... ] ) ]

Description

CREATE USER MAPPING defines a mapping of a user to a foreign server. A user mapping typically encapsulates connection information that a foreign-data wrapper uses together with the information encapsulated by a foreign server to access an external data resource.

The owner of a foreign server can create user mappings for that server for any user. Also, a user can create a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters

IF NOT EXISTS

Do not throw an error if a mapping of the given user to the given foreign server already exists. A notice is issued in this case. Note that there is no guarantee that the existing user mapping is anything like the one that would have been created.

user_name

The name of an existing user that is mapped to foreign server. CURRENT_USER and USER match the name of the current user. When PUBLIC is specified, a so-called public mapping is created that is used when no user-specific mapping is applicable.

server_name

The name of an existing server for which the user mapping is to be created.

OPTIONS ( option 'value' [ , ... ] )

This clause specifies the options of the user mapping. The options typically define the actual user name and password of the mapping. Option names must be unique. The allowed option names and values are specific to the server's foreign-data wrapper.

Examples

Create a user mapping for user bob, server foo:

CREATE USER MAPPING FOR bob SERVER foo OPTIONS (user 'bob', password 'secret');

Compatibility

CREATE USER MAPPING conforms to ISO/IEC 9075-9 (SQL/MED).
CREATE USER MAPPING

See Also
ALTER USER MAPPING, DROP USER MAPPING, CREATE FOREIGN DATA WRAPPER, CREATE SERVER
CREATE VIEW

CREATE VIEW — define a new view

Synopsis

```
CREATE [ OR REPLACE ] [ TEMP | TEMPORARY ] [ RECURSIVE ] VIEW name
[ ( column_name [, ...] ) ]
[ WITH ( view_option_name = view_option_value [ , ... ] ) ]
AS query
[ WITH [ CASCADED | LOCAL ] CHECK OPTION ]
```

Description

CREATE VIEW defines a view of a query. The view is not physically materialized. Instead, the query is run every time the view is referenced in a query.

CREATE OR REPLACE VIEW is similar, but if a view of the same name already exists, it is replaced. The new query must generate the same columns that were generated by the existing view query (that is, the same column names in the same order and with the same data types), but it may add additional columns to the end of the list. The calculations giving rise to the output columns may be completely different.

If a schema name is given (for example, CREATE VIEW myschema.myview ...) then the view is created in the specified schema. Otherwise it is created in the current schema. Temporary views exist in a special schema, so a schema name cannot be given when creating a temporary view. The name of the view must be distinct from the name of any other view, table, sequence, index or foreign table in the same schema.

Parameters

TEMPORARY or TEMP

If specified, the view is created as a temporary view. Temporary views are automatically dropped at the end of the current session. Existing permanent relations with the same name are not visible to the current session while the temporary view exists, unless they are referenced with schema-qualified names.

If any of the tables referenced by the view are temporary, the view is created as a temporary view (whether TEMPORARY is specified or not).

RECURSIVE

Creates a recursive view. The syntax

```
CREATE RECURSIVE VIEW [ schema . ] view_name (column_names) AS
SELECT ...;
```

is equivalent to

```
CREATE VIEW [ schema . ] view_name AS WITH RECURSIVE view_name
(column_names) AS (SELECT ...) SELECT column_names
FROM view_name;
```
CREATE VIEW

A view column name list must be specified for a recursive view.

name

The name (optionally schema-qualified) of a view to be created.

column_name

An optional list of names to be used for columns of the view. If not given, the column names are deduced from the query.

WITH ( view_option_name [= view_option_value] [, ... ] )

This clause specifies optional parameters for a view; the following parameters are supported:

check_option(string)

This parameter may be either local or cascaded, and is equivalent to specifying WITH [ CASCADED | LOCAL ] CHECK OPTION (see below). This option can be changed on existing views using ALTER VIEW.

security_barrier(boolean)

This should be used if the view is intended to provide row-level security. See Section 41.5 for full details.

query

A SELECT or VALUES command which will provide the columns and rows of the view.

WITH [ CASCADED | LOCAL ] CHECK OPTION

This option controls the behavior of automatically updatable views. When this option is specified, INSERT and UPDATE commands on the view will be checked to ensure that new rows satisfy the view-defining condition (that is, the new rows are checked to ensure that they are visible through the view). If they are not, the update will be rejected. If the CHECK OPTION is not specified, INSERT and UPDATE commands on the view are allowed to create rows that are not visible through the view. The following check options are supported:

LOCAL

New rows are only checked against the conditions defined directly in the view itself. Any conditions defined on underlying base views are not checked (unless they also specify the CHECK OPTION).

CASCADED

New rows are checked against the conditions of the view and all underlying base views. If the CHECK OPTION is specified, and neither LOCAL nor CASCADED is specified, then CASCADED is assumed.

The CHECK OPTION may not be used with RECURSIVE views.

Note that the CHECK OPTION is only supported on views that are automatically updatable, and do not have INSTEAD OF triggers or INSTEAD rules. If an automatically updatable view is defined on top of a base view that has INSTEAD OF triggers, then the LOCAL CHECK OPTION may be used to check the conditions on the automatically updatable view, but the conditions on the base view with INSTEAD OF triggers will not be checked (a cascaded check option will not cascade down to a trigger-updatable view, and any check options defined directly on a trigger-updatable view will be ignored). If the view or any of its base relations has an INSTEAD rule that causes the INSERT or UPDATE command to be rewritten, then all check options will be ignored in the
rewritten query, including any checks from automatically updatable views defined on top of the relation with the INSTEAD rule.

Notes

Use the DROP VIEW statement to drop views.

Be careful that the names and types of the view's columns will be assigned the way you want. For example:

```
CREATE VIEW vista AS SELECT 'Hello World';
```

is bad form because the column name defaults to `column`; also, the column data type defaults to `text`, which might not be what you wanted. Better style for a string literal in a view's result is something like:

```
CREATE VIEW vista AS SELECT text 'Hello World' AS hello;
```

Access to tables referenced in the view is determined by permissions of the view owner. In some cases, this can be used to provide secure but restricted access to the underlying tables. However, not all views are secure against tampering; see Section 41.5 for details. Functions called in the view are treated the same as if they had been called directly from the query using the view. Therefore the user of a view must have permissions to call all functions used by the view.

When CREATE OR REPLACE VIEW is used on an existing view, only the view’s defining SELECT rule is changed. Other view properties, including ownership, permissions, and non-SELECT rules, remain unchanged. You must own the view to replace it (this includes being a member of the owning role).

Updatable Views

Simple views are automatically updatable: the system will allow INSERT, UPDATE and DELETE statements to be used on the view in the same way as on a regular table. A view is automatically updatable if it satisfies all of the following conditions:

- The view must have exactly one entry in its FROM list, which must be a table or another updatable view.
- The view definition must not contain WITH, DISTINCT, GROUP BY, HAVING, LIMIT, or OFFSET clauses at the top level.
- The view definition must not contain set operations (UNION, INTERSECT or EXCEPT) at the top level.
- The view’s select list must not contain any aggregates, window functions or set-returning functions.

An automatically updatable view may contain a mix of updatable and non-updatable columns. A column is updatable if it is a simple reference to an updatable column of the underlying base relation; otherwise the column is read-only, and an error will be raised if an INSERT or UPDATE statement attempts to assign a value to it.

If the view is automatically updatable the system will convert any INSERT, UPDATE or DELETE statement on the view into the corresponding statement on the underlying base relation. INSERT statements that have an ON CONFLICT UPDATE clause are fully supported.

If an automatically updatable view contains a WHERE condition, the condition restricts which rows of the base relation are available to be modified by UPDATE and DELETE statements on the view. However, an UPDATE is allowed to change a row so that it no longer satisfies the WHERE condition,
CREATE VIEW
and thus is no longer visible through the view. Similarly, an INSERT command can potentially insert base-relation rows that do not satisfy the WHERE condition and thus are not visible through the view (ON CONFLICT UPDATE may similarly affect an existing row not visible through the view). The CHECK OPTION may be used to prevent INSERT and UPDATE commands from creating such rows that are not visible through the view.

If an automatically updatable view is marked with the security_barrier property then all the view’s WHERE conditions (and any conditions using operators which are marked as LEAKPROOF) will always be evaluated before any conditions that a user of the view has added. See Section 41.5 for full details. Note that, due to this, rows which are not ultimately returned (because they do not pass the user’s WHERE conditions) may still end up being locked. EXPLAIN can be used to see which conditions are applied at the relation level (and therefore do not lock rows) and which are not.

A more complex view that does not satisfy all these conditions is read-only by default: the system will not allow an insert, update, or delete on the view. You can get the effect of an updatable view by creating INSTEAD OF triggers on the view, which must convert attempted inserts, etc. on the view into appropriate actions on other tables. For more information see CREATE TRIGGER. Another possibility is to create rules (see CREATE RULE), but in practice triggers are easier to understand and use correctly.

Note that the user performing the insert, update or delete on the view must have the corresponding insert, update or delete privilege on the view. In addition the view’s owner must have the relevant privileges on the underlying base relations, but the user performing the update does not need any permissions on the underlying base relations (see Section 41.5).

Examples

Create a view consisting of all comedy films:

```sql
CREATE VIEW comedies AS
SELECT *
FROM films
WHERE kind = 'Comedy';
```

This will create a view containing the columns that are in the film table at the time of view creation. Though * was used to create the view, columns added later to the table will not be part of the view.

Create a view with LOCAL CHECK OPTION:

```sql
CREATE VIEW universal_comedies AS
SELECT *
FROM comedies
WHERE classification = 'U'
WITH LOCAL CHECK OPTION;
```

This will create a view based on the comedies view, showing only films with kind = 'Comedy' and classification = 'U'. Any attempt to INSERT or UPDATE a row in the view will be rejected if the new row doesn’t have classification = 'U', but the film kind will not be checked.

Create a view with CASCADED CHECK OPTION:

```sql
CREATE VIEW pg_comedies AS
SELECT *
FROM comedies
WHERE classification = 'PG'
```
WITH CASCADED CHECK OPTION;

This will create a view that checks both the kind and classification of new rows.

Create a view with a mix of updatable and non-updatable columns:

CREATE VIEW comedies AS
    SELECT f.*,
            country_code_to_name(f.country_code) AS country,
            (SELECT avg(r.rating)
             FROM user_ratings r
             WHERE r.film_id = f.id) AS avg_rating
    FROM films f
    WHERE f.kind = 'Comedy';

This view will support INSERT, UPDATE and DELETE. All the columns from the films table will be updatable, whereas the computed columns country and avg_rating will be read-only.

Create a recursive view consisting of the numbers from 1 to 100:

CREATE RECURSIVE VIEW public.nums_1_100 (n) AS
    VALUES (1)
    UNION ALL
    SELECT n+1 FROM nums_1_100 WHERE n < 100;

Notice that although the recursive view’s name is schema-qualified in this CREATE, its internal self-reference is not schema-qualified. This is because the implicitly-created CTE’s name cannot be schema-qualified.

**Compatibility**

CREATE OR REPLACE VIEW is a PostgreSQL language extension. So is the concept of a temporary view. The WITH ( ...) clause is an extension as well.

**See Also**

ALTER VIEW, DROP VIEW, CREATE MATERIALIZED VIEW
DEALLOCATE

DEALLOCATE — deallocate a prepared statement

Synopsis

DEALLOCATE [ PREPARE ] { name | ALL }

Description

DEALLOCATE is used to deallocate a previously prepared SQL statement. If you do not explicitly
deallocate a prepared statement, it is deallocated when the session ends.

For more information on prepared statements, see PREPARE.

Parameters

PREPARE

This key word is ignored.

name

The name of the prepared statement to deallocate.

ALL

Dealocate all prepared statements.

Compatibility

The SQL standard includes a DEALLOCATE statement, but it is only for use in embedded SQL.

See Also

EXECUTE, PREPARE
**DECLARE**

DECLARE — define a cursor

**Synopsis**

DECLARE name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ]
CURSOR [ { WITH | WITHOUT } HOLD ] FOR query

**Description**

DECLARE allows a user to create cursors, which can be used to retrieve a small number of rows at a time out of a larger query. After the cursor is created, rows are fetched from it using FETCH.

**Note**

This page describes usage of cursors at the SQL command level. If you are trying to use cursors inside a PL/pgSQL function, the rules are different — see Section 43.7.

**Parameters**

*name*

The name of the cursor to be created.

*BINARY*

Causes the cursor to return data in binary rather than in text format.

*INSENSITIVE*

Indicates that data retrieved from the cursor should be unaffected by updates to the table(s) underlying the cursor that occur after the cursor is created. In PostgreSQL, this is the default behavior; so this key word has no effect and is only accepted for compatibility with the SQL standard.

*SCROLL*

**NO SCROLL**

SCROLL specifies that the cursor can be used to retrieve rows in a nonsequential fashion (e.g., backward). Depending upon the complexity of the query's execution plan, specifying SCROLL might impose a performance penalty on the query's execution time. **NO SCROLL** specifies that the cursor cannot be used to retrieve rows in a nonsequential fashion. The default is to allow scrolling in some cases; this is not the same as specifying SCROLL. See Notes for details.

*WITH HOLD*

*WITHOUT HOLD*

WITH HOLD specifies that the cursor can continue to be used after the transaction that created it successfully commits. WITHOUT HOLD specifies that the cursor cannot be used outside of the transaction that created it. If neither WITHOUT HOLD nor WITH HOLD is specified, **WITHOUT HOLD** is the default.

*query*

A SELECT or VALUES command which will provide the rows to be returned by the cursor.
The key words BINARY, INSENSITIVE, and SCROLL can appear in any order.

Notes

Normal cursors return data in text format, the same as a SELECT would produce. The BINARY option specifies that the cursor should return data in binary format. This reduces conversion effort for both the server and client, at the cost of more programmer effort to deal with platform-dependent binary data formats. As an example, if a query returns a value of one from an integer column, you would get a string of 1 with a default cursor, whereas with a binary cursor you would get a 4-byte field containing the internal representation of the value (in big-endian byte order).

Binary cursors should be used carefully. Many applications, including psql, are not prepared to handle binary cursors and expect data to come back in the text format.

Note

When the client application uses the “extended query” protocol to issue a FETCH command, the Bind protocol message specifies whether data is to be retrieved in text or binary format. This choice overrides the way that the cursor is defined. The concept of a binary cursor as such is thus obsolete when using extended query protocol — any cursor can be treated as either text or binary.

Unless WITH HOLD is specified, the cursor created by this command can only be used within the current transaction. Thus, DECLARE without WITH HOLD is useless outside a transaction block: the cursor would survive only to the completion of the statement. Therefore PostgreSQL reports an error if such a command is used outside a transaction block. Use BEGIN and COMMIT (or ROLLBACK) to define a transaction block.

If WITH HOLD is specified and the transaction that created the cursor successfully commits, the cursor can continue to be accessed by subsequent transactions in the same session. (But if the creating transaction is aborted, the cursor is removed.) A cursor created with WITH HOLD is closed when an explicit CLOSE command is issued on it, or the session ends. In the current implementation, the rows represented by a held cursor are copied into a temporary file or memory area so that they remain available for subsequent transactions.

WITH HOLD may not be specified when the query includes FOR UPDATE or FOR SHARE.

The SCROLL option should be specified when defining a cursor that will be used to fetch backwards. This is required by the SQL standard. However, for compatibility with earlier versions, PostgreSQL will allow backward fetches without SCROLL, if the cursor’s query plan is simple enough that no extra overhead is needed to support it. However, application developers are advised not to rely on using backward fetches from a cursor that has not been created with SCROLL. If NO SCROLL is specified, then backward fetches are disallowed in any case.

Backward fetches are also disallowed when the query includes FOR UPDATE or FOR SHARE; therefore SCROLL may not be specified in this case.

Caution

Scrollable and WITH HOLD cursors may give unexpected results if they invoke any volatile functions (see Section 38.7). When a previously fetched row is re-fetched, the functions might be re-executed, perhaps leading to results different from the first time. One workaround for such cases is to declare the cursor WITH HOLD and commit the transaction before reading any rows from it. This will force the entire output of the cursor to be materialized in temporary storage, so that volatile functions are executed exactly once for each row.
DECLARE

If the cursor's query includes `FOR UPDATE` or `FOR SHARE`, then returned rows are locked at the
time they are first fetched, in the same way as for a regular SELECT command with these options. In
addition, the returned rows will be the most up-to-date versions; therefore these options provide the
equivalent of what the SQL standard calls a "sensitive cursor". (Specifying `INSENSITIVE` together
with `FOR UPDATE` or `FOR SHARE` is an error.)

Caution

It is generally recommended to use `FOR UPDATE` if the cursor is intended to be used
with `UPDATE ... WHERE CURRENT OF` or `DELETE ... WHERE CURRENT
OF`. Using `FOR UPDATE` prevents other sessions from changing the rows between the
time they are fetched and the time they are updated. Without `FOR UPDATE`, a subse-
quent `WHERE CURRENT OF` command will have no effect if the row was changed
since the cursor was created.

Another reason to use `FOR UPDATE` is that without it, a subsequent `WHERE CUR-
RENT OF` might fail if the cursor query does not meet the SQL standard's rules for
being "simply updatable" (in particular, the cursor must reference just one table and
not use grouping or `ORDER BY`). Cursors that are not simply updatable might work,
or might not, depending on plan choice details; so in the worst case, an application
might work in testing and then fail in production. If `FOR UPDATE` is specified, the
cursor is guaranteed to be updatable.

The main reason not to use `FOR UPDATE` with `WHERE CURRENT OF` is if you
need the cursor to be scrollable, or to be insensitive to the subsequent updates (that is,
continue to show the old data). If this is a requirement, pay close heed to the caveats
shown above.

The SQL standard only makes provisions for cursors in embedded SQL. The PostgreSQL server does
not implement an `OPEN` statement for cursors; a cursor is considered to be open when it is declared.
However, ECPG, the embedded SQL preprocessor for PostgreSQL, supports the standard SQL cursor
conventions, including those involving `DECLARE` and `OPEN` statements.

You can see all available cursors by querying the `pg_cursors` system view.

Examples

To declare a cursor:

```
DECLARE liahona CURSOR FOR SELECT * FROM films;
```

See FETCH for more examples of cursor usage.

Compatibility

The SQL standard says that it is implementation-dependent whether cursors are sensitive to concurrent
updates of the underlying data by default. In PostgreSQL, cursors are insensitive by default, and can
be made sensitive by specifying `FOR UPDATE`. Other products may work differently.

The SQL standard allows cursors only in embedded SQL and in modules. PostgreSQL permits cursors
to be used interactively.

Binary cursors are a PostgreSQL extension.

See Also

`CLOSE`, `FETCH`, `MOVE`
DELETE

DELETE — delete rows of a table

Synopsis

[ WITH [ RECURSIVE ] with_query [, ... ] ]
DELETE FROM [ ONLY ] table_name [ * ] [ [ AS ] alias ]
   [ USING using_list ]
   [ WHERE condition | WHERE CURRENT OF cursor_name ]
   [ RETURNING * | output_expression [ [ AS ] output_name ]
 [, ... ] ]

Description

DELETE deletes rows that satisfy the WHERE clause from the specified table. If the WHERE clause is absent, the effect is to delete all rows in the table. The result is a valid, but empty table.

Tip

TRUNCATE provides a faster mechanism to remove all rows from a table.

There are two ways to delete rows in a table using information contained in other tables in the database: using sub-selects, or specifying additional tables in the USING clause. Which technique is more appropriate depends on the specific circumstances.

The optional RETURNING clause causes DELETE to compute and return value(s) based on each row actually deleted. Any expression using the table's columns, and/or columns of other tables mentioned in USING, can be computed. The syntax of the RETURNING list is identical to that of the output list of SELECT.

You must have the DELETE privilege on the table to delete from it, as well as the SELECT privilege for any table in the USING clause or whose values are read in the condition.

Parameters

with_query

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the DELETE query. See Section 7.8 and SELECT for details.

table_name

The name (optionally schema-qualified) of the table to delete rows from. If ONLY is specified before the table name, matching rows are deleted from the named table only. If ONLY is not specified, matching rows are also deleted from any tables inheriting from the named table. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

alias

A substitute name for the target table. When an alias is provided, it completely hides the actual name of the table. For example, given DELETE FROM foo AS f, the remainder of the DELETE statement must refer to this table as f not foo.
**using_list**

A list of table expressions, allowing columns from other tables to appear in the `WHERE` condition. This is similar to the list of tables that can be specified in the `FROM` Clause of a `SELECT` statement; for example, an alias for the table name can be specified. Do not repeat the target table in the `using_list`, unless you wish to set up a self-join.

**condition**

An expression that returns a value of type `boolean`. Only rows for which this expression returns `true` will be deleted.

**cursor_name**

The name of the cursor to use in a `WHERE CURRENT OF` condition. The row to be deleted is the one most recently fetched from this cursor. The cursor must be a non-grouping query on the `DELETE`'s target table. Note that `WHERE CURRENT OF` cannot be specified together with a Boolean condition. See DECLARE for more information about using cursors with `WHERE CURRENT OF`.

**output_expression**

An expression to be computed and returned by the `DELETE` command after each row is deleted. The expression can use any column names of the table named by `table_name` or table(s) listed in `USING`. Write `*` to return all columns.

**output_name**

A name to use for a returned column.

**Outputs**

On successful completion, a `DELETE` command returns a command tag of the form

```
DELETE count
```

The `count` is the number of rows deleted. Note that the number may be less than the number of rows that matched the `condition` when deletes were suppressed by a `BEFORE DELETE` trigger. If `count` is 0, no rows were deleted by the query (this is not considered an error).

If the `DELETE` command contains a `RETURNING` clause, the result will be similar to that of a `SELECT` statement containing the columns and values defined in the `RETURNING` list, computed over the row(s) deleted by the command.

**Notes**

PostgreSQL lets you reference columns of other tables in the `WHERE` condition by specifying the other tables in the `USING` clause. For example, to delete all films produced by a given producer, one can do:

```
DELETE FROM films USING producers
    WHERE producer_id = producers.id AND producers.name = 'foo';
```

What is essentially happening here is a join between `films` and `producers`, with all successfully joined `films` rows being marked for deletion. This syntax is not standard. A more standard way to do it is:

```
DELETE FROM films
```
WHERE producer_id IN (SELECT id FROM producers WHERE name = 'foo');

In some cases the join style is easier to write or faster to execute than the sub-select style.

Examples

Delete all films but musicals:

```sql
DELETE FROM films WHERE kind <> 'Musical';
```

Clear the table `films`:

```sql
DELETE FROM films;
```

Delete completed tasks, returning full details of the deleted rows:

```sql
DELETE FROM tasks WHERE status = 'DONE' RETURNING *;
```

Delete the row of `tasks` on which the cursor `c_tasks` is currently positioned:

```sql
DELETE FROM tasks WHERE CURRENT OF c_tasks;
```

Compatibility

This command conforms to the SQL standard, except that the `USING` and `RETURNING` clauses are PostgreSQL extensions, as is the ability to use `WITH` with `DELETE`.

See Also

TRUNCATE
DISCARD

DISCARD — discard session state

Synopsis

DISCARD ( ALL | PLANS | SEQUENCES | TEMPORARY | TEMP )

Description

DISCARD releases internal resources associated with a database session. This command is useful for partially or fully resetting the session's state. There are several subcommands to release different types of resources; the DISCARD ALL variant subsumes all the others, and also resets additional state.

Parameters

PLANS

Releases all cached query plans, forcing re-planning to occur the next time the associated prepared statement is used.

SEQUENCES

Discards all cached sequence-related state, including currval()/lastval() information and any preallocated sequence values that have not yet been returned by nextval(). (See CREATE SEQUENCE for a description of preallocated sequence values.)

TEMPORARY or TEMP

Drops all temporary tables created in the current session.

ALL

Releases all temporary resources associated with the current session and resets the session to its initial state. Currently, this has the same effect as executing the following sequence of statements:

SET SESSION AUTHORIZATION DEFAULT;
RESET ALL;
DEALLOCATE ALL;
CLOSE ALL;
UNLISTEN *;
SELECT pg_advisory_unlock_all();
DISCARD PLANS;
DISCARD SEQUENCES;
DISCARD TEMP;

Notes

DISCARD ALL cannot be executed inside a transaction block.

Compatibility

DISCARD is a PostgreSQL extension.
**DO**

DO — execute an anonymous code block

**Synopsis**

DO [ LANGUAGE lang_name ] code

**Description**

DO executes an anonymous code block, or in other words a transient anonymous function in a procedural language.

The code block is treated as though it were the body of a function with no parameters, returning `void`. It is parsed and executed a single time.

The optional `LANGUAGE` clause can be written either before or after the code block.

**Parameters**

`code`

The procedural language code to be executed. This must be specified as a string literal, just as in `CREATE FUNCTION`. Use of a dollar-quoted literal is recommended.

`lang_name`

The name of the procedural language the code is written in. If omitted, the default is `plpgsql`.

**Notes**

The procedural language to be used must already have been installed into the current database by means of `CREATE EXTENSION`. `plpgsql` is installed by default, but other languages are not.

The user must have `USAGE` privilege for the procedural language, or must be a superuser if the language is untrusted. This is the same privilege requirement as for creating a function in the language.

If `DO` is executed in a transaction block, then the procedure code cannot execute transaction control statements. Transaction control statements are only allowed if `DO` is executed in its own transaction.

**Examples**

Grant all privileges on all views in schema `public` to role `webuser`:

```sql
DO $$
DECLARE r record;
BEGIN
  FOR r IN SELECT table_schema, table_name FROM information_schema.tables
    WHERE table_type = 'VIEW' AND table_schema = 'public'
  LOOP
    EXECUTE 'GRANT ALL ON ' || quote_ident(r.table_schema) || '.
    EXECUTE ' TO webuser';
  END LOOP;
$$;
```
DO
END$$;

Compatibility

There is no DO statement in the SQL standard.

See Also
CREATE LANGUAGE
DROP ACCESS METHOD

DROP ACCESS METHOD — remove an access method

Synopsis

DROP ACCESS METHOD [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP ACCESS METHOD removes an existing access method. Only superusers can drop access methods.

Parameters

IF EXISTS

Do not throw an error if the access method does not exist. A notice is issued in this case.

name

The name of an existing access method.

CASCADE

Automatically drop objects that depend on the access method (such as operator classes, operator families, and indexes), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the access method if any objects depend on it. This is the default.

Examples

Drop the access method heptree:

DROP ACCESS METHOD heptree;

Compatibility

DROP ACCESS METHOD is a PostgreSQL extension.

See Also

CREATE ACCESS METHOD
DROP AGGREGATE

DROP AGGREGATE — remove an aggregate function

Synopsis

DROP AGGREGATE [ IF EXISTS ] name ( aggregate_signature ) [, ...]  
[ CASCADE | RESTRICT ]

where aggregate_signature is:

* | 
[ argmode ] [ argname ] argtype [ , ... ] | 
[ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ] 
[ argname ] argtype [ , ... ]

Description

DROP AGGREGATE removes an existing aggregate function. To execute this command the current user must be the owner of the aggregate function.

Parameters

IF EXISTS

Do not throw an error if the aggregate does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing aggregate function.

argmode

The mode of an argument: IN or VARIADIC. If omitted, the default is IN.

argname

The name of an argument. Note that DROP AGGREGATE does not actually pay any attention to argument names, since only the argument data types are needed to determine the aggregate function's identity.

argtype

An input data type on which the aggregate function operates. To reference a zero-argument aggregate function, write * in place of the list of argument specifications. To reference an ordered-set aggregate function, write ORDER BY between the direct and aggregated argument specifications.

CASCADE

Automatically drop objects that depend on the aggregate function (such as views using it), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the aggregate function if any objects depend on it. This is the default.
Notes

Alternative syntaxes for referencing ordered-set aggregates are described under ALTER AGGREGATE.

Examples

To remove the aggregate function myavg for type integer:

```
DROP AGGREGATE myavg(integer);
```

To remove the hypothetical-set aggregate function myrank, which takes an arbitrary list of ordering columns and a matching list of direct arguments:

```
DROP AGGREGATE myrank(VARIADIC "any" ORDER BY VARIADIC "any");
```

To remove multiple aggregate functions in one command:

```
DROP AGGREGATE myavg(integer), myavg(bigint);
```

Compatibility

There is no DROP AGGREGATE statement in the SQL standard.

See Also

ALTER AGGREGATE, CREATE AGGREGATE
DROP CAST

DROP CAST — remove a cast

Synopsis

DROP CAST [ IF EXISTS ] (source_type AS target_type) [ CASCADE | RESTRICT ]

Description

DROP CAST removes a previously defined cast.

To be able to drop a cast, you must own the source or the target data type. These are the same privileges that are required to create a cast.

Parameters

IF EXISTS

Do not throw an error if the cast does not exist. A notice is issued in this case.

source_type

The name of the source data type of the cast.

target_type

The name of the target data type of the cast.

CASCADE

RESTRICT

These key words do not have any effect, since there are no dependencies on casts.

Examples

To drop the cast from type text to type int:

DROP CAST (text AS int);

Compatibility

The DROP CAST command conforms to the SQL standard.

See Also

CREATE CAST
DROP COLLATION

DROP COLLATION — remove a collation

Synopsis

```
DROP COLLATION [ IF EXISTS ] name [ CASCADE | RESTRICT ]
```

Description

DROP COLLATION removes a previously defined collation. To be able to drop a collation, you must own the collation.

Parameters

**IF EXISTS**

Do not throw an error if the collation does not exist. A notice is issued in this case.

**name**

The name of the collation. The collation name can be schema-qualified.

**CASCADE**

Automatically drop objects that depend on the collation, and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the collation if any objects depend on it. This is the default.

Examples

To drop the collation named german:

```
DROP COLLATION german;
```

Compatibility

The DROP COLLATION command conforms to the SQL standard, apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER COLLATION, CREATE COLLATION
DROP CONVERSION

DROP CONVERSION — remove a conversion

Synopsis

```
DROP CONVERSION [ IF EXISTS ] name [ CASCADE | RESTRICT ]
```

Description

DROP CONVERSION removes a previously defined conversion. To be able to drop a conversion, you must own the conversion.

Parameters

IF EXISTS

Do not throw an error if the conversion does not exist. A notice is issued in this case.

name

The name of the conversion. The conversion name can be schema-qualified.

CASCADE

RESTRICT

These key words do not have any effect, since there are no dependencies on conversions.

Examples

To drop the conversion named myname:

```
DROP CONVERSION myname;
```

Compatibility

There is no DROP CONVERSION statement in the SQL standard, but a DROP TRANSLATION statement that goes along with the CREATE TRANSLATION statement that is similar to the CREATE CONVERSION statement in PostgreSQL.

See Also

ALTER CONVERSION, CREATE CONVERSION
DROP DATABASE

DROP DATABASE — remove a database

Synopsis

```
DROP DATABASE [ IF EXISTS ] name
```

Description

DROP DATABASE drops a database. It removes the catalog entries for the database and deletes the directory containing the data. It can only be executed by the database owner. Also, it cannot be executed while you or anyone else are connected to the target database. (Connect to postgres or any other database to issue this command.)

DROP DATABASE cannot be undone. Use it with care!

Parameters

`IF EXISTS`

Do not throw an error if the database does not exist. A notice is issued in this case.

`name`

The name of the database to remove.

Notes

DROP DATABASE cannot be executed inside a transaction block.

This command cannot be executed while connected to the target database. Thus, it might be more convenient to use the program dropdb instead, which is a wrapper around this command.

Compatibility

There is no DROP DATABASE statement in the SQL standard.

See Also

CREATE DATABASE
DROP DOMAIN

DROP DOMAIN — remove a domain

Synopsis

```
DROP DOMAIN [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]
```

Description

DROP DOMAIN removes a domain. Only the owner of a domain can remove it.

Parameters

**IF EXISTS**

Do not throw an error if the domain does not exist. A notice is issued in this case.

**name**

The name (optionally schema-qualified) of an existing domain.

**CASCADE**

Automatically drop objects that depend on the domain (such as table columns), and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the domain if any objects depend on it. This is the default.

Examples

To remove the domain `box`:

```
DROP DOMAIN box;
```

Compatibility

This command conforms to the SQL standard, except for the IF EXISTS option, which is a PostgreSQL extension.

See Also

CREATE DOMAIN, ALTER DOMAIN
DROP EVENT TRIGGER

DROP EVENT TRIGGER — remove an event trigger

Synopsis

DROP EVENT TRIGGER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP EVENT TRIGGER removes an existing event trigger. To execute this command, the current user must be the owner of the event trigger.

Parameters

IF EXISTS

Do not throw an error if the event trigger does not exist. A notice is issued in this case.

name

The name of the event trigger to remove.

CASCADE

Automatically drop objects that depend on the trigger, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the trigger if any objects depend on it. This is the default.

Examples

Destroy the trigger snitch:

DROP EVENT TRIGGER snitch;

Compatibility

There is no DROP EVENT TRIGGER statement in the SQL standard.

See Also

CREATE EVENT TRIGGER, ALTER EVENT TRIGGER
DROP EXTENSION

DROP EXTENSION — remove an extension

Synopsis

```
DROP EXTENSION [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]
```

Description

DROP EXTENSION removes extensions from the database. Dropping an extension causes its component objects to be dropped as well.

You must own the extension to use DROP EXTENSION.

Parameters

IF EXISTS

Do not throw an error if the extension does not exist. A notice is issued in this case.

name

The name of an installed extension.

CASCADE

Automatically drop objects that depend on the extension, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the extension if any objects depend on it (other than its own member objects and other extensions listed in the same DROP command). This is the default.

Examples

To remove the extension hstore from the current database:

```
DROP EXTENSION hstore;
```

This command will fail if any of hstore's objects are in use in the database, for example if any tables have columns of the hstore type. Add the CASCADE option to forcibly remove those dependent objects as well.

Compatibility

DROP EXTENSION is a PostgreSQL extension.

See Also

CREATE EXTENSION, ALTER EXTENSION
DROP FOREIGN DATA WRAPPER

DROP FOREIGN DATA WRAPPER — remove a foreign-data wrapper

Synopsis

DROP FOREIGN DATA WRAPPER [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP FOREIGN DATA WRAPPER removes an existing foreign-data wrapper. To execute this command, the current user must be the owner of the foreign-data wrapper.

Parameters

IF EXISTS

Do not throw an error if the foreign-data wrapper does not exist. A notice is issued in this case.

name

The name of an existing foreign-data wrapper.

CASCADE

Automatically drop objects that depend on the foreign-data wrapper (such as foreign tables and servers), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the foreign-data wrapper if any objects depend on it. This is the default.

Examples

Drop the foreign-data wrapper dbi:

DROP FOREIGN DATA WRAPPER dbi;

Compatibility

DROP FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.

See Also

CREATE FOREIGN DATA WRAPPER, ALTER FOREIGN DATA WRAPPER
DROP FOREIGN TABLE

DROP FOREIGN TABLE — remove a foreign table

Synopsis

DROP FOREIGN TABLE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP FOREIGN TABLE removes a foreign table. Only the owner of a foreign table can remove it.

Parameters

IF EXISTS

Do not throw an error if the foreign table does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the foreign table to drop.

CASCADE

Automatically drop objects that depend on the foreign table (such as views), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the foreign table if any objects depend on it. This is the default.

Examples

To destroy two foreign tables, films and distributors:

DROP FOREIGN TABLE films, distributors;

Compatibility

This command conforms to the ISO/IEC 9075-9 (SQL/MED), except that the standard only allows one foreign table to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER FOREIGN TABLE, CREATE FOREIGN TABLE
DROP FUNCTION

DROP FUNCTION — remove a function

Synopsis

DROP FUNCTION [ IF EXISTS ] name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) [, ...] ] [ CASCADE | RESTRICT ]

Description

DROP FUNCTION removes the definition of an existing function. To execute this command the user must be the owner of the function. The argument types to the function must be specified, since several different functions can exist with the same name and different argument lists.

Parameters

IF EXISTS

Do not throw an error if the function does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing function. If no argument list is specified, the name must be unique in its schema.

argmode

The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that DROP FUNCTION does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function's identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname

The name of an argument. Note that DROP FUNCTION does not actually pay any attention to argument names, since only the argument data types are needed to determine the function's identity.

argtype

The data type(s) of the function's arguments (optionally schema-qualified), if any.

CASCADE

Automatically drop objects that depend on the function (such as operators or triggers), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the function if any objects depend on it. This is the default.

Examples

This command removes the square root function:
DROP FUNCTION sqrt(integer);

Drop multiple functions in one command:

DROP FUNCTION sqrt(integer), sqrt(bigint);

If the function name is unique in its schema, it can be referred to without an argument list:

DROP FUNCTION update_employee_salaries;

Note that this is different from

DROP FUNCTION update_employee_salaries();

which refers to a function with zero arguments, whereas the first variant can refer to a function with any number of arguments, including zero, as long as the name is unique.

Compatibility

This command conforms to the SQL standard, with these PostgreSQL extensions:

• The standard only allows one function to be dropped per command.

• The IF EXISTS option

• The ability to specify argument modes and names

See Also

CREATE FUNCTION, ALTER FUNCTION, DROP PROCEDURE, DROP ROUTINE
DROP GROUP

DROP GROUP — remove a database role

Synopsis

```
DROP GROUP [ IF EXISTS ] name [, ...]
```

Description

DROP GROUP is now an alias for DROP ROLE.

Compatibility

There is no DROP GROUP statement in the SQL standard.

See Also

DROP ROLE
DROP INDEX

DROP INDEX — remove an index

Synopsis

```sql
DROP INDEX [ CONCURRENTLY ] [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]
```

Description

DROP INDEX drops an existing index from the database system. To execute this command you must be the owner of the index.

Parameters

CONCURRENTLY

Drop the index without locking out concurrent selects, inserts, updates, and deletes on the index’s table. A normal DROP INDEX acquires exclusive lock on the table, blocking other accesses until the index drop can be completed. With this option, the command instead waits until conflicting transactions have completed.

There are several caveats to be aware of when using this option. Only one index name can be specified, and the CASCADE option is not supported. (Thus, an index that supports a UNIQUE or PRIMARY KEY constraint cannot be dropped this way.) Also, regular DROP INDEX commands can be performed within a transaction block, but DROP INDEX CONCURRENTLY cannot.

IF EXISTS

Do not throw an error if the index does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an index to remove.

CASCADE

Automatically drop objects that depend on the index, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the index if any objects depend on it. This is the default.

Examples

This command will remove the index title_idx:

```sql
DROP INDEX title_idx;
```

Compatibility

DROP INDEX is a PostgreSQL language extension. There are no provisions for indexes in the SQL standard.
See Also
CREATE INDEX
DROP LANGUAGE

DROP LANGUAGE — remove a procedural language

Synopsis

DROP [ PROCEDURAL ] LANGUAGE [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP LANGUAGE removes the definition of a previously registered procedural language. You must be a superuser or the owner of the language to use DROP LANGUAGE.

Note

As of PostgreSQL 9.1, most procedural languages have been made into “extensions”, and should therefore be removed with DROP EXTENSION not DROP LANGUAGE.

Parameters

IF EXISTS

Do not throw an error if the language does not exist. A notice is issued in this case.

name

The name of an existing procedural language. For backward compatibility, the name can be enclosed by single quotes.

CASCADE

Automatically drop objects that depend on the language (such as functions in the language), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the language if any objects depend on it. This is the default.

Examples

This command removes the procedural language plsample:

DROP LANGUAGE plsample;

Compatibility

There is no DROP LANGUAGE statement in the SQL standard.

See Also

ALTER LANGUAGE, CREATE LANGUAGE
DROP MATERIALIZED VIEW

DROP MATERIALIZED VIEW — remove a materialized view

Synopsis

DROP MATERIALIZED VIEW [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP MATERIALIZED VIEW drops an existing materialized view. To execute this command you must be the owner of the materialized view.

Parameters

IF EXISTS

Do not throw an error if the materialized view does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the materialized view to remove.

CASCADE

Automatically drop objects that depend on the materialized view (such as other materialized views, or regular views), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the materialized view if any objects depend on it. This is the default.

Examples

This command will remove the materialized view called order_summary:

DROP MATERIALIZED VIEW order_summary;

Compatibility

DROP MATERIALIZED VIEW is a PostgreSQL extension.

See Also

CREATE MATERIALIZED VIEW, ALTER MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
DROP OPERATOR

DROP OPERATOR — remove an operator

Synopsis

DROP OPERATOR [ IF EXISTS ] name ( { left_type | NONE } ,
( right_type | NONE } ) [, ...] [ CASCADE | RESTRICT ]

Description

DROP OPERATOR drops an existing operator from the database system. To execute this command you must be the owner of the operator.

Parameters

IF EXISTS

Do not throw an error if the operator does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing operator.

left_type

The data type of the operator's left operand; write NONE if the operator has no left operand.

right_type

The data type of the operator's right operand; write NONE if the operator has no right operand.

CASCADE

Automatically drop objects that depend on the operator (such as views using it), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the operator if any objects depend on it. This is the default.

Examples

Remove the power operator \( a^b \) for type integer:

DROP OPERATOR ^ (integer, integer);

Remove the left unary bitwise complement operator \( \sim b \) for type bit:

DROP OPERATOR \( \sim \) (none, bit);

Remove the right unary factorial operator \( x! \) for type bigint:

1667
DROP OPERATOR ! (bigint, none);

Remove multiple operators in one command:

DROP OPERATOR ~ (none, bit), ! (bigint, none);

**Compatibility**

There is no DROP OPERATOR statement in the SQL standard.

**See Also**

CREATE OPERATOR, ALTER OPERATOR
DROP OPERATOR CLASS

DROP OPERATOR CLASS — remove an operator class

Synopsis

DROP OPERATOR CLASS [ IF EXISTS ] name USING index_method [ CASCADE | RESTRICT ]

Description

DROP OPERATOR CLASS drops an existing operator class. To execute this command you must be
the owner of the operator class.

DROP OPERATOR CLASS does not drop any of the operators or functions referenced by the class.
If there are any indexes depending on the operator class, you will need to specify CASCADE for the
drop to complete.

Parameters

IF EXISTS
Do not throw an error if the operator class does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of an existing operator class.

index_method
The name of the index access method the operator class is for.

CASCADE
Automatically drop objects that depend on the operator class (such as indexes), and in turn all
objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the operator class if any objects depend on it. This is the default.

Notes

DROP OPERATOR CLASS will not drop the operator family containing the class, even if there is
nothing else left in the family (in particular, in the case where the family was implicitly created by
CREATE OPERATOR CLASS). An empty operator family is harmless, but for the sake of tidiness
you might wish to remove the family with DROP OPERATOR FAMILY; or perhaps better, use DROP
OPERATOR FAMILY in the first place.

Examples

Remove the B-tree operator class widget_ops:

DROP OPERATOR CLASS widget_ops USING btree;
This command will not succeed if there are any existing indexes that use the operator class. Add CASCADE to drop such indexes along with the operator class.

Compatibility

There is no DROP OPERATOR CLASS statement in the SQL standard.

See Also

ALTER OPERATOR CLASS, CREATE OPERATOR CLASS, DROP OPERATOR FAMILY
DROP OPERATOR FAMILY

DROP OPERATOR FAMILY — remove an operator family

Synopsis

```
DROP OPERATOR FAMILY [ IF EXISTS ] name USING index_method
[ CASCADE | RESTRICT ]
```

Description

DROP OPERATOR FAMILY drops an existing operator family. To execute this command you must
be the owner of the operator family.

DROP OPERATOR FAMILY includes dropping any operator classes contained in the family, but
it does not drop any of the operators or functions referenced by the family. If there are any indexes
depending on operator classes within the family, you will need to specify CASCADE for the drop to
complete.

Parameters

IF EXISTS

Do not throw an error if the operator family does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing operator family.

index_method

The name of the index access method the operator family is for.

CASCADE

Automatically drop objects that depend on the operator family, and in turn all objects that depend
on those objects (see Section 5.13).

RESTRICT

Refuse to drop the operator family if any objects depend on it. This is the default.

Examples

Remove the B-tree operator family float_ops:

```
DROP OPERATOR FAMILY float_ops USING btree;
```

This command will not succeed if there are any existing indexes that use operator classes within the
family. Add CASCADE to drop such indexes along with the operator family.

Compatibility

There is no DROP OPERATOR FAMILY statement in the SQL standard.
See Also

ALTER OPERATOR FAMILY, CREATE OPERATOR FAMILY, ALTER OPERATOR CLASS,
CREATE OPERATOR CLASS, DROP OPERATOR CLASS
DROP OWNED

DROP OWNED — remove database objects owned by a database role

Synopsis

DROP OWNED BY { name | CURRENT_USER | SESSION_USER } [, ...] [ CASCADE | RESTRICT ]

Description

DROP OWNED drops all the objects within the current database that are owned by one of the specified roles. Any privileges granted to the given roles on objects in the current database and on shared objects (databases, tablespaces) will also be revoked.

Parameters

name

The name of a role whose objects will be dropped, and whose privileges will be revoked.

CASCADE

Automatically drop objects that depend on the affected objects, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the objects owned by a role if any other database objects depend on one of the affected objects. This is the default.

Notes

DROP OWNED is often used to prepare for the removal of one or more roles. Because DROP OWNED only affects the objects in the current database, it is usually necessary to execute this command in each database that contains objects owned by a role that is to be removed.

Using the CASCADE option might make the command recurse to objects owned by other users.

The REASSIGN OWNED command is an alternative that reassigns the ownership of all the database objects owned by one or more roles. However, REASSIGN OWNED does not deal with privileges for other objects.

Databases and tablespaces owned by the role(s) will not be removed.

See Section 21.4 for more discussion.

Compatibility

The DROP OWNED command is a PostgreSQL extension.

See Also

REASSIGN OWNED, DROP ROLE
DROP POLICY

DROP POLICY — remove a row level security policy from a table

Synopsis

DROP POLICY [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

Description

DROP POLICY removes the specified policy from the table. Note that if the last policy is removed for a table and the table still has row level security enabled via ALTER TABLE, then the default-deny policy will be used. ALTER TABLE ... DISABLE ROW LEVEL SECURITY can be used to disable row level security for a table, whether policies for the table exist or not.

Parameters

IF EXISTS

Do not throw an error if the policy does not exist. A notice is issued in this case.

name

The name of the policy to drop.

table_name

The name (optionally schema-qualified) of the table that the policy is on.

CASCADE
RESTRICT

These key words do not have any effect, since there are no dependencies on policies.

Examples

To drop the policy called p1 on the table named my_table:

DROP POLICY p1 ON my_table;

Compatibility

DROP POLICY is a PostgreSQL extension.

See Also

CREATE POLICY, ALTER POLICY
DROP PROCEDURE

DROP PROCEDURE — remove a procedure

Synopsis

DROP PROCEDURE [ IF EXISTS ] name [ ( [ argmode ] [ argname ] argtype [, ...] ) [, ...] ]

[ CASCADE | RESTRICT ]

Description

DROP PROCEDURE removes the definition of an existing procedure. To execute this command the
user must be the owner of the procedure. The argument types to the procedure must be specified, since
several different procedures can exist with the same name and different argument lists.

Parameters

IF EXISTS

Do not throw an error if the procedure does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing procedure. If no argument list is specified,
the name must be unique in its schema.

argmode

The mode of an argument: IN or VARIADIC. If omitted, the default is IN.

argname

The name of an argument. Note that DROP PROCEDURE does not actually pay any attention
to argument names, since only the argument data types are needed to determine the procedure's
identity.

argtype

The data type(s) of the procedure's arguments (optionally schema-qualified), if any.

CASCADE

Automatically drop objects that depend on the procedure, and in turn all objects that depend on
those objects (see Section 5.13).

RESTRICT

Refuse to drop the procedure if any objects depend on it. This is the default.

Examples

DROP PROCEDURE do_db_maintenance();
Compatibility

This command conforms to the SQL standard, with these PostgreSQL extensions:

- The standard only allows one procedure to be dropped per command.
- The `IF EXISTS` option
- The ability to specify argument modes and names

See Also

CREATE PROCEDURE, ALTER PROCEDURE, DROP FUNCTION, DROP ROUTINE
DROP PUBLICATION

DROP PUBLICATION — remove a publication

Synopsis

DROP PUBLICATION [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP PUBLICATION removes an existing publication from the database.

A publication can only be dropped by its owner or a superuser.

Parameters

IF EXISTS

Do not throw an error if the publication does not exist. A notice is issued in this case.

name

The name of an existing publication.

CASCADE
RESTRIC

These key words do not have any effect, since there are no dependencies on publications.

Examples

Drop a publication:

DROP PUBLICATION mypublication;

Compatibility

DROP PUBLICATION is a PostgreSQL extension.

See Also

CREATE PUBLICATION, ALTER PUBLICATION
DROP ROLE

DROP ROLE — remove a database role

Synopsis

DROP ROLE [ IF EXISTS ] name [, ...]

Description

DROP ROLE removes the specified role(s). To drop a superuser role, you must be a superuser yourself; to drop non-superuser roles, you must have CREATEROLE privilege.

A role cannot be removed if it is still referenced in any database of the cluster; an error will be raised if so. Before dropping the role, you must drop all the objects it owns (or reassign their ownership) and revoke any privileges the role has been granted on other objects. The REASSIGN OWNED and DROP OWNED commands can be useful for this purpose; see Section 21.4 for more discussion.

However, it is not necessary to remove role memberships involving the role; DROP ROLE automatically revokes any memberships of the target role in other roles, and of other roles in the target role. The other roles are not dropped nor otherwise affected.

Parameters

IF EXISTS

Do not throw an error if the role does not exist. A notice is issued in this case.

name

The name of the role to remove.

Notes

PostgreSQL includes a program dropuser that has the same functionality as this command (in fact, it calls this command) but can be run from the command shell.

Examples

To drop a role:

DROP ROLE jonathan;

Compatibility

The SQL standard defines DROP ROLE, but it allows only one role to be dropped at a time, and it specifies different privilege requirements than PostgreSQL uses.

See Also

CREATE ROLE, ALTER ROLE, SET ROLE
DROP ROUTINE

DROP ROUTINE — remove a routine

Synopsis

DROP ROUTINE [ IF EXISTS ] name [ ( [ [ argmode ] [ argname ] argtype [, ... ] ) ] [, ...] [ CASCADE | RESTRICT ]

Description

DROP ROUTINE removes the definition of an existing routine, which can be an aggregate function, a normal function, or a procedure. See under DROP AGGREGATE, DROP FUNCTION, and DROP PROCEDURE for the description of the parameters, more examples, and further details.

Examples

To drop the routine foo for type integer:

DROP ROUTINE foo(integer);

This command will work independent of whether foo is an aggregate, function, or procedure.

Compatibility

This command conforms to the SQL standard, with these PostgreSQL extensions:

• The standard only allows one routine to be dropped per command.
• The IF EXISTS option
• The ability to specify argument modes and names
• Aggregate functions are an extension.

See Also

DROP AGGREGATE, DROP FUNCTION, DROP PROCEDURE, ALTER ROUTINE

Note that there is no CREATE ROUTINE command.
DROP RULE

-- remove a rewrite rule

Synopsis

DROP RULE [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

Description

DROP RULE drops a rewrite rule.

Parameters

IF EXISTS

Do not throw an error if the rule does not exist. A notice is issued in this case.

name

The name of the rule to drop.

table_name

The name (optionally schema-qualified) of the table or view that the rule applies to.

CASCADE

Automatically drop objects that depend on the rule, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the rule if any objects depend on it. This is the default.

Examples

To drop the rewrite rule newrule:

DROP RULE newrule ON mytable;

Compatibility

DROP RULE is a PostgreSQL language extension, as is the entire query rewrite system.

See Also

CREATE RULE, ALTER RULE
DROP SCHEMA

DROP SCHEMA — remove a schema

Synopsis

DROP SCHEMA [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP SCHEMA removes schemas from the database.

A schema can only be dropped by its owner or a superuser. Note that the owner can drop the schema (and thereby all contained objects) even if they do not own some of the objects within the schema.

Parameters

IF EXISTS

Do not throw an error if the schema does not exist. A notice is issued in this case.

name

The name of a schema.

CASCADE

Automatically drop objects (tables, functions, etc.) that are contained in the schema, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the schema if it contains any objects. This is the default.

Notes

Using the CASCADE option might make the command remove objects in other schemas besides the one(s) named.

Examples

To remove schema mystuff from the database, along with everything it contains:

DROP SCHEMA mystuff CASCADE;

Compatibility

DROP SCHEMA is fully conforming with the SQL standard, except that the standard only allows one schema to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER SCHEMA, CREATE SCHEMA
DROP SEQUENCE

DROP SEQUENCE — remove a sequence

Synopsis

```
DROP SEQUENCE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]
```

Description

DROP SEQUENCE removes sequence number generators. A sequence can only be dropped by its owner or a superuser.

Parameters

IF EXISTS

Do not throw an error if the sequence does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of a sequence.

CASCADE

Automatically drop objects that depend on the sequence, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the sequence if any objects depend on it. This is the default.

Examples

To remove the sequence serial:

```
DROP SEQUENCE serial;
```

Compatibility

DROP SEQUENCE conforms to the SQL standard, except that the standard only allows one sequence to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

CREATE SEQUENCE, ALTER SEQUENCE
DROP SERVER

DROP SERVER — remove a foreign server descriptor

Synopsis

DROP SERVER [ IF EXISTS ] name [, ... ] [ CASCADE | RESTRICT ]

Description

DROP SERVER removes an existing foreign server descriptor. To execute this command, the current user must be the owner of the server.

Parameters

IF EXISTS

Do not throw an error if the server does not exist. A notice is issued in this case.

name

The name of an existing server.

CASCADE

Automatically drop objects that depend on the server (such as user mappings), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the server if any objects depend on it. This is the default.

Examples

Drop a server foo if it exists:

DROP SERVER IF EXISTS foo;

Compatibility

DROP SERVER conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.

See Also

CREATE SERVER, ALTER SERVER
DROP STATISTICS

Synopsis

DROP STATISTICS [ IF EXISTS ] name [, ...]

Description

DROP STATISTICS removes statistics object(s) from the database. Only the statistics object's owner, the schema owner, or a superuser can drop a statistics object.

Parameters

IF EXISTS

Do not throw an error if the statistics object does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the statistics object to drop.

Examples

To destroy two statistics objects in different schemas, without failing if they don't exist:

DROP STATISTICS IF EXISTS
    accounting.users_uid_creation,
    public.grants_user_role;

Compatibility

There is no DROP STATISTICS command in the SQL standard.

See Also

ALTER STATISTICS, CREATE STATISTICS
DROP SUBSCRIPTION

DROP SUBSCRIPTION — remove a subscription

Synopsis

```
DROP SUBSCRIPTION [ IF EXISTS ] name [ CASCADE | RESTRICT ]
```

Description

DROP SUBSCRIPTION removes a subscription from the database cluster.

A subscription can only be dropped by a superuser.

DROP SUBSCRIPTION cannot be executed inside a transaction block if the subscription is associated with a replication slot. (You can use ALTER SUBSCRIPTION to unset the slot.)

Parameters

- `name`

  The name of a subscription to be dropped.

- `CASCADE`

- `RESTRICT`

  These key words do not have any effect, since there are no dependencies on subscriptions.

Notes

When dropping a subscription that is associated with a replication slot on the remote host (the normal state), DROP SUBSCRIPTION will connect to the remote host and try to drop the replication slot as part of its operation. This is necessary so that the resources allocated for the subscription on the remote host are released. If this fails, either because the remote host is not reachable or because the remote replication slot cannot be dropped or does not exist or never existed, the DROP SUBSCRIPTION command will fail. To proceed in this situation, disassociate the subscription from the replication slot by executing ALTER SUBSCRIPTION ... SET (slot_name = NONE). After that, DROP SUBSCRIPTION will no longer attempt any actions on a remote host. Note that if the remote replication slot still exists, it should then be dropped manually; otherwise it will continue to reserve WAL and might eventually cause the disk to fill up. See also Section 31.2.1.

If a subscription is associated with a replication slot, then DROP SUBSCRIPTION cannot be executed inside a transaction block.

Examples

Drop a subscription:

```
DROP SUBSCRIPTION mysub;
```

Compatibility

DROP SUBSCRIPTION is a PostgreSQL extension.
See Also

CREATE SUBSCRIPTION, ALTER SUBSCRIPTION
DROP TABLE

DROP TABLE — remove a table

Synopsis

DROP TABLE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP TABLE removes tables from the database. Only the table owner, the schema owner, and superuser can drop a table. To empty a table of rows without destroying the table, use DELETE or TRUNCATE.

DROP TABLE always removes any indexes, rules, triggers, and constraints that exist for the target table. However, to drop a table that is referenced by a view or a foreign-key constraint of another table, CASCADE must be specified. (CASCADE will remove a dependent view entirely, but in the foreign-key case it will only remove the foreign-key constraint, not the other table entirely.)

Parameters

IF EXISTS

Do not throw an error if the table does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the table to drop.

CASCADE

Automatically drop objects that depend on the table (such as views), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the table if any objects depend on it. This is the default.

Examples

To destroy two tables, films and distributors:

DROP TABLE films, distributors;

Compatibility

This command conforms to the SQL standard, except that the standard only allows one table to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER TABLE, CREATE TABLE
DROP TABLESPACE

DROP TABLESPACE — remove a tablespace

Synopsis

DROP TABLESPACE [ IF EXISTS ] name

Description

DROP TABLESPACE removes a tablespace from the system.

A tablespace can only be dropped by its owner or a superuser. The tablespace must be empty of all database objects before it can be dropped. It is possible that objects in other databases might still reside in the tablespace even if no objects in the current database are using the tablespace. Also, if the tablespace is listed in the temp_tablespaces setting of any active session, the DROP might fail due to temporary files residing in the tablespace.

Parameters

IF EXISTS

Do not throw an error if the tablespace does not exist. A notice is issued in this case.

name

The name of a tablespace.

Notes

DROP TABLESPACE cannot be executed inside a transaction block.

Examples

To remove tablespace mystuff from the system:

DROP TABLESPACE mystuff;

Compatibility

DROP TABLESPACE is a PostgreSQL extension.

See Also

CREATE TABLESPACE, ALTER TABLESPACE
DROP TEXT SEARCH CONFIGURATION

DROP TEXT SEARCH CONFIGURATION — remove a text search configuration

Synopsis

DROP TEXT SEARCH CONFIGURATION [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP TEXT SEARCH CONFIGURATION drops an existing text search configuration. To execute this command you must be the owner of the configuration.

Parameters

IF EXISTS

Do not throw an error if the text search configuration does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing text search configuration.

CASCADE

Automatically drop objects that depend on the text search configuration, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the text search configuration if any objects depend on it. This is the default.

Examples

Remove the text search configuration my_english:

DROP TEXT SEARCH CONFIGURATION my_english;

This command will not succeed if there are any existing indexes that reference the configuration in to_tsvector calls. Add CASCADE to drop such indexes along with the text search configuration.

Compatibility

There is no DROP TEXT SEARCH CONFIGURATION statement in the SQL standard.

See Also

ALTER TEXT SEARCH CONFIGURATION, CREATE TEXT SEARCH CONFIGURATION
DROP TEXT SEARCH DICTIONARY

DROP TEXT SEARCH DICTIONARY — remove a text search dictionary

Synopsis

```
DROP TEXT SEARCH DICTIONARY [ IF EXISTS ] name [ CASCADE |
RESTRICT ]
```

Description

DROP TEXT SEARCH DICTIONARY drops an existing text search dictionary. To execute this command you must be the owner of the dictionary.

Parameters

**IF EXISTS**

Do not throw an error if the text search dictionary does not exist. A notice is issued in this case.

**name**

The name (optionally schema-qualified) of an existing text search dictionary.

**CASCADE**

Automatically drop objects that depend on the text search dictionary, and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the text search dictionary if any objects depend on it. This is the default.

Examples

Remove the text search dictionary `english`:

```
DROP TEXT SEARCH DICTIONARY english;
```

This command will not succeed if there are any existing text search configurations that use the dictionary. Add `CASCADE` to drop such configurations along with the dictionary.

Compatibility

There is no `DROP TEXT SEARCH DICTIONARY` statement in the SQL standard.

See Also

ALTER TEXT SEARCH DICTIONARY, CREATE TEXT SEARCH DICTIONARY
DROP TEXT SEARCH PARSER

DROP TEXT SEARCH PARSER — remove a text search parser

Synopsis

DROP TEXT SEARCH PARSER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP TEXT SEARCH PARSER drops an existing text search parser. You must be a superuser to use this command.

Parameters

IF EXISTS

Do not throw an error if the text search parser does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing text search parser.

CASCADE

Automatically drop objects that depend on the text search parser, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the text search parser if any objects depend on it. This is the default.

Examples

Remove the text search parser my_parser:

DROP TEXT SEARCH PARSER my_parser;

This command will not succeed if there are any existing text search configurations that use the parser. Add CASCADE to drop such configurations along with the parser.

Compatibility

There is no DROP TEXT SEARCH PARSER statement in the SQL standard.

See Also

ALTER TEXT SEARCH PARSER, CREATE TEXT SEARCH PARSER
DROP TEXT SEARCH TEMPLATE

DROP TEXT SEARCH TEMPLATE — remove a text search template

Synopsis

DROP TEXT SEARCH TEMPLATE [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP TEXT SEARCH TEMPLATE drops an existing text search template. You must be a superuser to use this command.

Parameters

IF EXISTS

Do not throw an error if the text search template does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing text search template.

CASCADE

Automatically drop objects that depend on the text search template, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the text search template if any objects depend on it. This is the default.

Examples

Remove the text search template thesaurus:

DROP TEXT SEARCH TEMPLATE thesaurus;

This command will not succeed if there are any existing text search dictionaries that use the template. Add CASCADE to drop such dictionaries along with the template.

Compatibility

There is no DROP TEXT SEARCH TEMPLATE statement in the SQL standard.

See Also

ALTER TEXT SEARCH TEMPLATE, CREATE TEXT SEARCH TEMPLATE
DROP TRANSFORM

DROP TRANSFORM — remove a transform

Synopsis

```
DROP TRANSFORM [ IF EXISTS ] FOR type_name LANGUAGE lang_name
[ CASCADE | RESTRICT ]
```

Description

DROP TRANSFORM removes a previously defined transform.

To be able to drop a transform, you must own the type and the language. These are the same privileges that are required to create a transform.

Parameters

IF EXISTS

Do not throw an error if the transform does not exist. A notice is issued in this case.

```
type_name
```

The name of the data type of the transform.

```
lang_name
```

The name of the language of the transform.

CASCADE

Automatically drop objects that depend on the transform, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the transform if any objects depend on it. This is the default.

Examples

To drop the transform for type hstore and language plpythonu:

```
DROP TRANSFORM FOR hstore LANGUAGE plpythonu;
```

Compatibility

This form of DROP TRANSFORM is a PostgreSQL extension. See CREATE TRANSFORM for details.

See Also

CREATE TRANSFORM
DROP TRIGGER

DROP TRIGGER — remove a trigger

Synopsis

DROP TRIGGER [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

Description

DROP TRIGGER removes an existing trigger definition. To execute this command, the current user must be the owner of the table for which the trigger is defined.

Parameters

IF EXISTS

Do not throw an error if the trigger does not exist. A notice is issued in this case.

name

The name of the trigger to remove.

table_name

The name (optionally schema-qualified) of the table for which the trigger is defined.

CASCADE

Automatically drop objects that depend on the trigger, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the trigger if any objects depend on it. This is the default.

Examples

Destroy the trigger if_dist_exists on the table films:

DROP TRIGGER if_dist_exists ON films;

Compatibility

The DROP TRIGGER statement in PostgreSQL is incompatible with the SQL standard. In the SQL standard, trigger names are not local to tables, so the command is simply DROP TRIGGER name.

See Also

CREATE TRIGGER
DROP TYPE

DROP TYPE — remove a data type

Synopsis

DROP TYPE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP TYPE removes a user-defined data type. Only the owner of a type can remove it.

Parameters

IF EXISTS

Do not throw an error if the type does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the data type to remove.

CASCADE

Automatically drop objects that depend on the type (such as table columns, functions, and operators), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the type if any objects depend on it. This is the default.

Examples

To remove the data type box:

DROP TYPE box;

Compatibility

This command is similar to the corresponding command in the SQL standard, apart from the IF EXISTS option, which is a PostgreSQL extension. But note that much of the CREATE TYPE command and the data type extension mechanisms in PostgreSQL differ from the SQL standard.

See Also

ALTER TYPE, CREATE TYPE
DROP USER

DROP USER — remove a database role

Synopsis

```
DROP USER [ IF EXISTS ] name [, ...]
```

Description

DROP USER is simply an alternate spelling of DROP ROLE.

Compatibility

The DROP USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also

DROP ROLE
DROP USER MAPPING

DROP USER MAPPING — remove a user mapping for a foreign server

Synopsis

\[
\text{DROP USER MAPPING \{ IF EXISTS \} FOR \{ user_name | USER | CURRENT_USER | PUBLIC \} SERVER server_name}
\]

Description

DROP USER MAPPING removes an existing user mapping from foreign server.

The owner of a foreign server can drop user mappings for that server for any user. Also, a user can drop a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters

\begin{itemize}
\item \text{IF EXISTS}
\end{itemize}

Do not throw an error if the user mapping does not exist. A notice is issued in this case.

\begin{itemize}
\item \text{user_name}
\end{itemize}

User name of the mapping. \text{CURRENT_USER} and \text{USER} match the name of the current user. \text{PUBLIC} is used to match all present and future user names in the system.

\begin{itemize}
\item \text{server_name}
\end{itemize}

Server name of the user mapping.

Examples

Drop a user mapping \text{bob}, server \text{foo} if it exists:

\[
\text{DROP USER MAPPING IF EXISTS FOR bob SERVER foo;}
\]

Compatibility

\[
\text{DROP USER MAPPING} \text{ conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.}
\]

See Also

CREATE USER MAPPING, ALTER USER MAPPING
DROP VIEW

DROP VIEW — remove a view

Synopsis

```
DROP VIEW [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]
```

Description

DROP VIEW drops an existing view. To execute this command you must be the owner of the view.

Parameters

**IF EXISTS**

Do not throw an error if the view does not exist. A notice is issued in this case.

**name**

The name (optionally schema-qualified) of the view to remove.

**CASCADE**

Automatically drop objects that depend on the view (such as other views), and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the view if any objects depend on it. This is the default.

Examples

This command will remove the view called kinds:

```
DROP VIEW kinds;
```

Compatibility

This command conforms to the SQL standard, except that the standard only allows one view to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER VIEW, CREATE VIEW
**END**

END — commit the current transaction

**Synopsis**

END [ WORK | TRANSACTION ]

**Description**

END commits the current transaction. All changes made by the transaction become visible to others and are guaranteed to be durable if a crash occurs. This command is a PostgreSQL extension that is equivalent to COMMIT.

**Parameters**

WORK
TRANSACTION

Optional key words. They have no effect.

**Notes**

Use ROLLBACK to abort a transaction.

Issuing END when not inside a transaction does no harm, but it will provoke a warning message.

**Examples**

To commit the current transaction and make all changes permanent:

END;

**Compatibility**

END is a PostgreSQL extension that provides functionality equivalent to COMMIT, which is specified in the SQL standard.

**See Also**

BEGIN, COMMIT, ROLLBACK
EXECUTE

EXECUTE — execute a prepared statement

Synopsis

EXECUTE name [ ( parameter [, ...] ) ]

Description

EXECUTE is used to execute a previously prepared statement. Since prepared statements only exist for the duration of a session, the prepared statement must have been created by a PREPARE statement executed earlier in the current session.

If the PREPARE statement that created the statement specified some parameters, a compatible set of parameters must be passed to the EXECUTE statement, or else an error is raised. Note that (unlike functions) prepared statements are not overloaded based on the type or number of their parameters; the name of a prepared statement must be unique within a database session.

For more information on the creation and usage of prepared statements, see PREPARE.

Parameters

name

The name of the prepared statement to execute.

parameter

The actual value of a parameter to the prepared statement. This must be an expression yielding a value that is compatible with the data type of this parameter, as was determined when the prepared statement was created.

Outputs

The command tag returned by EXECUTE is that of the prepared statement, and not EXECUTE.

Examples

Examples are given in the Examples section of the PREPARE documentation.

Compatibility

The SQL standard includes an EXECUTE statement, but it is only for use in embedded SQL. This version of the EXECUTE statement also uses a somewhat different syntax.

See Also

DEALLOCATE, PREPARE
**EXPLAIN**

EXPLAIN — show the execution plan of a statement

**Synopsis**

```
EXPLAIN [ ( option [, ...] ) ] statement
EXPLAIN [ ANALYZE ] [ VERBOSE ] statement
```

where *option* can be one of:

- ANALYZE [ boolean ]
- VERBOSE [ boolean ]
- COSTS [ boolean ]
- BUFFERS [ boolean ]
- TIMING [ boolean ]
- SUMMARY [ boolean ]
- FORMAT { TEXT | XML | JSON | YAML }

**Description**

This command displays the execution plan that the PostgreSQL planner generates for the supplied statement. The execution plan shows how the table(s) referenced by the statement will be scanned — by plain sequential scan, index scan, etc. — and if multiple tables are referenced, what join algorithms will be used to bring together the required rows from each input table.

The most critical part of the display is the estimated statement execution cost, which is the planner's guess at how long it will take to run the statement (measured in cost units that are arbitrary, but conventionally mean disk page fetches). Actually two numbers are shown: the start-up cost before the first row can be returned, and the total cost to return all the rows. For most queries the total cost is what matters, but in contexts such as a subquery in EXISTS, the planner will choose the smallest start-up cost instead of the smallest total cost (since the executor will stop after getting one row, anyway). Also, if you limit the number of rows to return with a LIMIT clause, the planner makes an appropriate interpolation between the endpoint costs to estimate which plan is really the cheapest.

The ANALYZE option causes the statement to be actually executed, not only planned. Then actual run time statistics are added to the display, including the total elapsed time expended within each plan node (in milliseconds) and the total number of rows it actually returned. This is useful for seeing whether the planner's estimates are close to reality.

**Important**

Keep in mind that the statement is actually executed when the ANALYZE option is used. Although EXPLAIN will discard any output that a SELECT would return, other side effects of the statement will happen as usual. If you wish to use EXPLAIN ANALYZE on an INSERT, UPDATE, DELETE, CREATE TABLE AS, or EXECUTE statement without letting the command affect your data, use this approach:

```
BEGIN;
EXPLAIN ANALYZE ...;
ROLLBACK;
```
Only the ANALYZE and VERBOSE options can be specified, and only in that order, without surrounding the option list in parentheses. Prior to PostgreSQL 9.0, the unparenthesized syntax was the only one supported. It is expected that all new options will be supported only in the parenthesized syntax.

**Parameters**

**ANALYZE**

Carry out the command and show actual run times and other statistics. This parameter defaults to FALSE.

**VERBOSE**

Display additional information regarding the plan. Specifically, include the output column list for each node in the plan tree, schema-qualify table and function names, always label variables in expressions with their range table alias, and always print the name of each trigger for which statistics are displayed. This parameter defaults to FALSE.

**COSTS**

Include information on the estimated startup and total cost of each plan node, as well as the estimated number of rows and the estimated width of each row. This parameter defaults to TRUE.

**BUFFERS**

Include information on buffer usage. Specifically, include the number of shared blocks hit, read, dirtied, and written, the number of local blocks hit, read, dirtied, and written, and the number of temp blocks read and written. A hit means that a read was avoided because the block was found already in cache when needed. Shared blocks contain data from regular tables and indexes; local blocks contain data from temporary tables and indexes; while temp blocks contain short-term working data used in sorts, hashes, Materialize plan nodes, and similar cases. The number of blocks dirtied indicates the number of previously unmodified blocks that were changed by this query; while the number of blocks written indicates the number of previously-dirtied blocks evicted from cache by this backend during query processing. The number of blocks shown for an upper-level node includes those used by all its child nodes. In text format, only non-zero values are printed. This parameter may only be used when ANALYZE is also enabled. It defaults to FALSE.

**TIMING**

Include actual startup time and time spent in each node in the output. The overhead of repeatedly reading the system clock can slow down the query significantly on some systems, so it may be useful to set this parameter to FALSE when only actual row counts, and not exact times, are needed. Run time of the entire statement is always measured, even when node-level timing is turned off with this option. This parameter may only be used when ANALYZE is also enabled. It defaults to TRUE.

**SUMMARY**

Include summary information (e.g., totaled timing information) after the query plan. Summary information is included by default when ANALYZE is used but otherwise is not included by default, but can be enabled using this option. Planning time in EXPLAIN EXECUTE includes the time required to fetch the plan from the cache and the time required for re-planning, if necessary.

**FORMAT**

Specify the output format, which can be TEXT, XML, JSON, or YAML. Non-text output contains the same information as the text output format, but is easier for programs to parse. This parameter defaults to TEXT.
**boolean**

Specifies whether the selected option should be turned on or off. You can write `TRUE`, `ON`, or `1` to enable the option, and `FALSE`, `OFF`, or `0` to disable it. The `boolean` value can also be omitted, in which case `TRUE` is assumed.

**statement**

Any `SELECT`, `INSERT`, `UPDATE`, `DELETE`, `VALUES`, `EXECUTE`, `DECLARE`, `CREATE TABLE AS`, or `CREATE MATERIALIZED VIEW AS` statement, whose execution plan you wish to see.

**Outputs**

The command's result is a textual description of the plan selected for the `statement`, optionally annotated with execution statistics. Section 14.1 describes the information provided.

**Notes**

In order to allow the PostgreSQL query planner to make reasonably informed decisions when optimizing queries, the `pg_statistic` data should be up-to-date for all tables used in the query. Normally the autovacuum daemon will take care of that automatically. But if a table has recently had substantial changes in its contents, you might need to do a manual `ANALYZE` rather than wait for autovacuum to catch up with the changes.

In order to measure the run-time cost of each node in the execution plan, the current implementation of `EXPLAIN ANALYZE` adds profiling overhead to query execution. As a result, running `EXPLAIN ANALYZE` on a query can sometimes take significantly longer than executing the query normally. The amount of overhead depends on the nature of the query, as well as the platform being used. The worst case occurs for plan nodes that in themselves require very little time per execution, and on machines that have relatively slow operating system calls for obtaining the time of day.

**Examples**

To show the plan for a simple query on a table with a single `integer` column and 10000 rows:

```sql
EXPLAIN SELECT * FROM foo;
```

```
QUERY PLAN
---------------------------------------------------------
Seq Scan on foo  (cost=0.00..155.00 rows=10000 width=4)
(1 row)
```

Here is the same query, with JSON output formatting:

```sql
EXPLAIN (FORMAT JSON) SELECT * FROM foo;
```

```
| [ | + |
| } | + |
| "Plan": { | + |
| "Node Type": "Seq Scan", | + |
| "Relation Name": "foo", | + |
| "Alias": "foo", | + |
| "Startup Cost": 0.00, | + |
| "Total Cost": 155.00, | + |
| "Plan Rows": 10000, | + |
```
If there is an index and we use a query with an indexable WHERE condition, EXPLAIN might show a different plan:

```sql
EXPLAIN SELECT * FROM foo WHERE i = 4;
```

**QUERY PLAN**

```
Index Scan using fi on foo  (cost=0.00..5.98 rows=1 width=4)
Index Cond: (i = 4)
```

Here is the same query, but in YAML format:

```sql
EXPLAIN (FORMAT YAML) SELECT * FROM foo WHERE i='4';
```

**QUERY PLAN**

```
- Plan:                      +
  Node Type: "Index Scan"  +
  Scan Direction: "Forward"+  
  Index Name: "fi"  +
  Relation Name: "foo"  +
  Alias: "foo"  +
  Startup Cost: 0.00       +
  Total Cost: 5.98         +
  Plan Rows: 1             +
  Plan Width: 4            +
  Index Cond: "(i = 4)"
```

XML format is left as an exercise for the reader.

Here is the same plan with cost estimates suppressed:

```sql
EXPLAIN (COSTS FALSE) SELECT * FROM foo WHERE i = 4;
```

**QUERY PLAN**

```
Index Scan using fi on foo
Index Cond: (i = 4)
```

Here is an example of a query plan for a query using an aggregate function:

```sql
EXPLAIN SELECT sum(i) FROM foo WHERE i < 10;
```

**QUERY PLAN**

```
Aggregate  (cost=23.93..23.93 rows=1 width=4)
  ->  Index Scan using fi on foo  (cost=0.00..23.92 rows=6 width=4)
```
Here is an example of using `EXPLAIN EXECUTE` to display the execution plan for a prepared query:

```sql
PREPARE query(int, int) AS SELECT sum(bar) FROM test
    WHERE id > $1 AND id < $2
    GROUP BY foo;

EXPLAIN ANALYZE EXECUTE query(100, 200);
```

**QUERY PLAN**

```
HashAggregate  (cost=9.54..9.54 rows=1 width=8) (actual time=0.156..0.161 rows=11 loops=1)
  Group Key: foo
    ->  Index Scan using test_pkey on test  (cost=0.29..9.29 rows=50 width=8) (actual time=0.039..0.091 rows=99 loops=1)
        Index Cond: ((id > $1) AND (id < $2))
Planning time: 0.197 ms
Execution time: 0.225 ms
```

Of course, the specific numbers shown here depend on the actual contents of the tables involved. Also note that the numbers, and even the selected query strategy, might vary between PostgreSQL releases due to planner improvements. In addition, the `ANALYZE` command uses random sampling to estimate data statistics; therefore, it is possible for cost estimates to change after a fresh run of `ANALYZE`, even if the actual distribution of data in the table has not changed.

### Compatibility

There is no `EXPLAIN` statement defined in the SQL standard.

### See Also

`ANALYZE`
FETCH

FETCH — retrieve rows from a query using a cursor

Synopsis

FETCH [ direction [ FROM | IN ] ] cursor_name

where direction can be empty or one of:

NEXT
PRIOR
FIRST
LAST
ABSOLUTE count
RELATIVE count
count
ALL
FORWARD
FORWARD count
FORWARD ALL
BACKWARD
BACKWARD count
BACKWARD ALL

Description

FETCH retrieves rows using a previously-created cursor.

A cursor has an associated position, which is used by FETCH. The cursor position can be before the first row of the query result, on any particular row of the result, or after the last row of the result. When created, a cursor is positioned before the first row. After fetching some rows, the cursor is positioned on the row most recently retrieved. If FETCH runs off the end of the available rows then the cursor is left positioned after the last row, or before the first row if fetching backward. FETCH ALL or FETCH BACKWARD ALL will always leave the cursor positioned after the last row or before the first row.

The forms NEXT, PRIOR, FIRST, LAST, ABSOLUTE, RELATIVE fetch a single row after moving the cursor appropriately. If there is no such row, an empty result is returned, and the cursor is left positioned before the first row or after the last row as appropriate.

The forms using FORWARD and BACKWARD retrieve the indicated number of rows moving in the forward or backward direction, leaving the cursor positioned on the last-returned row (or after/before all rows, if the count exceeds the number of rows available).

RELATIVE 0, FORWARD 0, and BACKWARD 0 all request fetching the current row without moving the cursor, that is, re-fetching the most recently fetched row. This will succeed unless the cursor is positioned before the first row or after the last row; in which case, no row is returned.

Note

This page describes usage of cursors at the SQL command level. If you are trying to use cursors inside a PL/pgSQL function, the rules are different — see Section 43.7.3.
Parameters

direction

direction defines the fetch direction and number of rows to fetch. It can be one of the following:

NEXT

Fetch the next row. This is the default if direction is omitted.

PRIOR

Fetch the prior row.

FIRST

Fetch the first row of the query (same as ABSOLUTE 1).

LAST

Fetch the last row of the query (same as ABSOLUTE -1).

ABSOLUTE count

Fetch the count'th row of the query, or the abs(count)'th row from the end if count is negative. Position before first row or after last row if count is out of range; in particular, ABSOLUTE 0 positions before the first row.

RELATIVE count

Fetch the count'th succeeding row, or the abs(count)'th prior row if count is negative. RELATIVE 0 re-fetches the current row, if any.

count

Fetch the next count rows (same as FORWARD count).

ALL

Fetch all remaining rows (same as FORWARD ALL).

FORWARD

Fetch the next row (same as NEXT).

FORWARD count

Fetch the next count rows. FORWARD 0 re-fetches the current row.

FORWARD ALL

Fetch all remaining rows.

BACKWARD

Fetch the prior row (same as PRIOR).

BACKWARD count

Fetch the prior count rows (scanning backwards). BACKWARD 0 re-fetches the current row.
**FETCH**

**BACKWARD ALL**

Fetch all prior rows (scanning backwards).

**count**

*count* is a possibly-signed integer constant, determining the location or number of rows to fetch. For **FORWARD** and **BACKWARD** cases, specifying a negative *count* is equivalent to changing the sense of **FORWARD** and **BACKWARD**.

**cursor_name**

An open cursor’s name.

**Outputs**

On successful completion, a **FETCH** command returns a command tag of the form

`FETCH count`

The *count* is the number of rows fetched (possibly zero). Note that in *psql*, the command tag will not actually be displayed, since *psql* displays the fetched rows instead.

**Notes**

The cursor should be declared with the **SCROLL** option if one intends to use any variants of **FETCH** other than **FETCH NEXT** or **FETCH FORWARD** with a positive count. For simple queries *PostgreSQL* will allow backwards fetch from cursors not declared with **SCROLL**, but this behavior is best not relied on. If the cursor is declared with **NO SCROLL**, no backward fetches are allowed.

**ABSOLUTE** fetches are not any faster than navigating to the desired row with a relative move: the underlying implementation must traverse all the intermediate rows anyway. Negative absolute fetches are even worse: the query must be read to the end to find the last row, and then traversed backward from there. However, rewinding to the start of the query (as with **FETCH ABSOLUTE 0**) is fast.

**DECLARE** is used to define a cursor. Use **MOVE** to change cursor position without retrieving data.

**Examples**

The following example traverses a table using a cursor:

```
BEGIN WORK;

-- Set up a cursor:
DECLARE liahona SCROLL CURSOR FOR SELECT * FROM films;

-- Fetch the first 5 rows in the cursor liahona:
FETCH FORWARD 5 FROM liahona;

| code     |          title          | did | date_prod  |   kind   |
|----------|-------------------------+-----+------------+----------|
| BL101    | The Third Man           | 101 | 1949-12-23 | Drama    |
| 01:44    |                         |     |            |          |
| BL102    | The African Queen       | 101 | 1951-08-11 | Romantic |
| 01:43    |                         |     |            |          |
```
Compatibility

The SQL standard defines \texttt{FETCH} for use in embedded SQL only. The variant of \texttt{FETCH} described here returns the data as if it were a \texttt{SELECT} result rather than placing it in host variables. Other than this point, \texttt{FETCH} is fully upward-compatible with the SQL standard.

The \texttt{FETCH} forms involving \texttt{FORWARD} and \texttt{BACKWARD}, as well as the forms \texttt{FETCH \ count} and \texttt{FETCH ALL}, in which \texttt{FORWARD} is implicit, are PostgreSQL extensions.

The SQL standard allows only \texttt{FROM} preceding the cursor name; the option to use \texttt{IN}, or to leave them out altogether, is an extension.

See Also

\texttt{CLOSE}, \texttt{DECLARE}, \texttt{MOVE}
GRANT

GRANT — define access privileges

Synopsis

GRANT { { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES
| TRIGGER }
[, ...] | ALL [ PRIVILEGES ] }
ON { [ TABLE ] table_name [, ...]
| ALL TABLES IN SCHEMA schema_name [, ...] }
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { { SELECT | INSERT | UPDATE | REFERENCES }
( column_name [, ...] )
[, ...] | ALL [ PRIVILEGES ] ( column_name [, ...] ) }
ON [ TABLE ] table_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { ( USAGE | SELECT | UPDATE )
[, ...] | ALL [ PRIVILEGES ] }
ON [ SEQUENCE sequence_name [, ...]
| ALL SEQUENCES IN SCHEMA schema_name [, ...] }
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { ( CREATE | CONNECT | TEMPORARY | TEMP ) [, ...] | ALL
[ PRIVILEGES ] }
ON DATABASE database_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON DOMAIN domain_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON FOREIGN DATA WRAPPER fdw_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON FOREIGN SERVER server_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { EXECUTE | ALL [ PRIVILEGES ] }
ON { { FUNCTION | PROCEDURE | ROUTINE } routine_name
( [ [ argmode ] [ arg_name ] arg_type [, ...] ] ) [, ...]
| ALL { FUNCTIONS | PROCEDURES | ROUTINES } IN
SCHEMA schema_name [, ...] }
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON LANGUAGE lang_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { { SELECT | UPDATE } [, ...] | ALL [ PRIVILEGES ] }
GRANT

ON LARGE OBJECT loid [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { ( CREATE | USAGE ) [, ...] | ALL [ PRIVILEGES ] }
ON SCHEMA schema_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { CREATE | ALL [ PRIVILEGES ] }
ON TABLESPACE tablespace_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON TYPE type_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

where role_specification can be:

[ GROUP ] role_name
| PUBLIC
| CURRENT_USER
| SESSION_USER

GRANT role_name [, ...] TO role_name [, ...] [ WITH ADMIN OPTION ]

Description

The GRANT command has two basic variants: one that grants privileges on a database object (table, column, view, foreign table, sequence, database, foreign-data wrapper, foreign server, function, procedure, procedural language, schema, or tablespace), and one that grants membership in a role. These variants are similar in many ways, but they are different enough to be described separately.

GRANT on Database Objects

This variant of the GRANT command gives specific privileges on a database object to one or more roles. These privileges are added to those already granted, if any.

There is also an option to grant privileges on all objects of the same type within one or more schemas. This functionality is currently supported only for tables, sequences, functions, and procedures. ALL TABLES also affects views and foreign tables, just like the specific-object GRANT command. ALL FUNCTIONS also affects aggregate functions, but not procedures, again just like the specific-object GRANT command.

The key word PUBLIC indicates that the privileges are to be granted to all roles, including those that might be created later. PUBLIC can be thought of as an implicitly defined group that always includes all roles. Any particular role will have the sum of privileges granted directly to it, privileges granted to any role it is presently a member of, and privileges granted to PUBLIC.

If WITH GRANT OPTION is specified, the recipient of the privilege can in turn grant it to others. Without a grant option, the recipient cannot do that. Grant options cannot be granted to PUBLIC.

There is no need to grant privileges to the owner of an object (usually the user that created it), as the owner has all privileges by default. (The owner could, however, choose to revoke some of their own privileges for safety.)

The right to drop an object, or to alter its definition in any way, is not treated as a grantable privilege: it is inherent in the owner, and cannot be granted or revoked. (However, a similar effect can be obtained by granting or revoking membership in the role that owns the object; see below.) The owner implicitly has all grant options for the object, too.
PostgreSQL grants default privileges on some types of objects to PUBLIC. No privileges are granted to PUBLIC by default on tables, table columns, sequences, foreign data wrappers, foreign servers, large objects, schemas, or tablespaces. For other types of objects, the default privileges granted to PUBLIC are as follows: CONNECT and TEMPORARY (create temporary tables) privileges for databases; EXECUTE privilege for functions and procedures; and USAGE privilege for languages and data types (including domains). The object owner can, of course, REVOKE both default and expressly granted privileges. (For maximum security, issue the REVOKE in the same transaction that creates the object; then there is no window in which another user can use the object.) Also, these initial default privilege settings can be changed using the ALTER DEFAULT PRIVILEGES command.

The possible privileges are:

**SELECT**

Allows SELECT from any column, or the specific columns listed, of the specified table, view, or sequence. Also allows the use of COPY TO. This privilege is also needed to reference existing column values in UPDATE or DELETE. For sequences, this privilege also allows the use of the currval function. For large objects, this privilege allows the object to be read.

**INSERT**

Allows INSERT of a new row into the specified table. If specific columns are listed, only those columns may be assigned to in the INSERT command (other columns will therefore receive default values). Also allows COPY FROM.

**UPDATE**

Allows UPDATE of any column, or the specific columns listed, of the specified table. (In practice, any nontrivial UPDATE command will require SELECT privilege as well, since it must reference table columns to determine which rows to update, and/or to compute new values for columns.) SELECT ... FOR UPDATE and SELECT ... FOR SHARE also require this privilege on at least one column, in addition to the SELECT privilege. For sequences, this privilege allows the use of the nextval and setval functions. For large objects, this privilege allows writing or truncating the object.

**DELETE**

Allows DELETE of a row from the specified table. (In practice, any nontrivial DELETE command will require SELECT privilege as well, since it must reference table columns to determine which rows to delete.)

**TRUNCATE**

Allows TRUNCATE on the specified table.

**REFERENCES**

Allows creation of a foreign key constraint referencing the specified table, or specified column(s) of the table. (See the CREATE TABLE statement.)

**TRIGGER**

Allows the creation of a trigger on the specified table. (See the CREATE TRIGGER statement.)

**CREATE**

For databases, allows new schemas and publications to be created within the database. For schemas, allows new objects to be created within the schema. To rename an existing object, you must own the object and have this privilege for the containing schema.
For tablespaces, allows tables, indexes, and temporary files to be created within the tablespace, and allows databases to be created that have the tablespace as their default tablespace. (Note that revoking this privilege will not alter the placement of existing objects.)

CONNECT

Allows the user to connect to the specified database. This privilege is checked at connection startup (in addition to checking any restrictions imposed by `pg_hba.conf`).

TEMPORARY

Allows temporary tables to be created while using the specified database.

EXECUTE

Allows the use of the specified function or procedure and the use of any operators that are implemented on top of the function. This is the only type of privilege that is applicable to functions and procedures. The `FUNCTION` syntax also works for aggregate functions. Alternatively, use `ROUTINE` to refer to a function, aggregate function, or procedure regardless of what it is.

USAGE

For procedural languages, allows the use of the specified language for the creation of functions in that language. This is the only type of privilege that is applicable to procedural languages.

For schemas, allows access to objects contained in the specified schema (assuming that the objects' own privilege requirements are also met). Essentially this allows the grantee to “look up” objects within the schema. Without this permission, it is still possible to see the object names, e.g. by querying the system tables. Also, after revoking this permission, existing backends might have statements that have previously performed this lookup, so this is not a completely secure way to prevent object access.

For sequences, this privilege allows the use of the `currval` and `nextval` functions.

For types and domains, this privilege allows the use of the type or domain in the creation of tables, functions, and other schema objects. (Note that it does not control general “usage” of the type, such as values of the type appearing in queries. It only prevents objects from being created that depend on the type. The main purpose of the privilege is controlling which users create dependencies on a type, which could prevent the owner from changing the type later.)

For foreign-data wrappers, this privilege allows creation of new servers using the foreign-data wrapper.

For servers, this privilege allows creation of foreign tables using the server. Grantees may also create, alter, or drop their own user mappings associated with that server.

ALL PRIVILEGES

Grant all of the available privileges at once. The `PRIVILEGES` key word is optional in PostgreSQL, though it is required by strict SQL.

The privileges required by other commands are listed on the reference page of the respective command.

**GRANT on Roles**

This variant of the `GRANT` command grants membership in a role to one or more other roles. Membership in a role is significant because it conveys the privileges granted to a role to each of its members.

If `WITH ADMIN OPTION` is specified, the member can in turn grant membership in the role to others, and revoke membership in the role as well. Without the admin option, ordinary users cannot do that. A
role is not considered to hold WITH ADMIN OPTION on itself, but it may grant or revoke membership in itself from a database session where the session user matches the role. Database superusers can grant or revoke membership in any role to anyone. Roles having CREATEROLE privilege can grant or revoke membership in any role that is not a superuser.

Unlike the case with privileges, membership in a role cannot be granted to PUBLIC. Note also that this form of the command does not allow the noise word GROUP.

Notes

The REVOKE command is used to revoke access privileges.

Since PostgreSQL 8.1, the concepts of users and groups have been unified into a single kind of entity called a role. It is therefore no longer necessary to use the keyword GROUP to identify whether a grantee is a user or a group. GROUP is still allowed in the command, but it is a noise word.

A user may perform SELECT, INSERT, etc. on a column if they hold that privilege for either the specific column or its whole table. Granting the privilege at the table level and then revoking it for one column will not do what one might wish: the table-level grant is unaffected by a column-level operation.

When a non-owner of an object attempts to GRANT privileges on the object, the command will fail outright if the user has no privileges whatsoever on the object. As long as some privilege is available, the command will proceed, but it will grant only those privileges for which the user has grant options. The GRANT ALL PRIVILEGES forms will issue a warning message if no grant options are held, while the other forms will issue a warning if grant options for any of the privileges specifically named in the command are not held. (In principle these statements apply to the object owner as well, but since the owner is always treated as holding all grant options, the cases can never occur.)

It should be noted that database superusers can access all objects regardless of object privilege settings. This is comparable to the rights of root in a Unix system. As with root, it's unwise to operate as a superuser except when absolutely necessary.

If a superuser chooses to issue a GRANT or REVOKE command, the command is performed as though it were issued by the owner of the affected object. In particular, privileges granted via such a command will appear to have been granted by the object owner. (For role membership, the membership appears to have been granted by the containing role itself.)

GRANT and REVOKE can also be done by a role that is not the owner of the affected object, but is a member of the role that owns the object, or is a member of a role that holds privileges WITH GRANT OPTION on the object. In this case the privileges will be recorded as having been granted by the role that actually owns the object or holds the privileges WITH GRANT OPTION. For example, if table t1 is owned by role g1, of which role u1 is a member, then u1 can grant privileges on t1 to u2, but those privileges will appear to have been granted directly by g1. Any other member of role g1 could revoke them later.

If the role executing GRANT holds the required privileges indirectly via more than one role membership path, it is unspecified which containing role will be recorded as having done the grant. In such cases it is best practice to use SET ROLE to become the specific role you want to do the GRANT as.

Granting permission on a table does not automatically extend permissions to any sequences used by the table, including sequences tied to SERIAL columns. Permissions on sequences must be set separately.

Use psql's \dp command to obtain information about existing privileges for tables and columns. For example:

```sql
=> \dp mytable
Access privileges
```
## GRANT

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Type</th>
<th>Access privileges</th>
<th>Column access privileges</th>
</tr>
</thead>
<tbody>
<tr>
<td>public</td>
<td>mytable</td>
<td>table</td>
<td>miriam=arwdDxt/miriam</td>
<td>col1: r/miriam, miriam_rw=rw/</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>miriam</td>
<td>admin=arw/miriam</td>
</tr>
</tbody>
</table>

(1 row)

The entries shown by `\dp` are interpreted thus:

- `rolename=xxxx` -- privileges granted to a role
- `=xxxx` -- privileges granted to PUBLIC
- `r` -- SELECT ("read")
- `w` -- UPDATE ("write")
- `a` -- INSERT ("append")
- `d` -- DELETE
- `D` -- TRUNCATE
- `x` -- REFERENCES
- `t` -- TRIGGER
- `X` -- EXECUTE
- `U` -- USAGE
- `C` -- CREATE
- `c` -- CONNECT
- `T` -- TEMPORARY

- `arwdDxt` -- ALL PRIVILEGES (for tables, varies for other objects)
- `*` -- grant option for preceding privilege
- `/yyyy` -- role that granted this privilege

The above example display would be seen by user `miriam` after creating table `mytable` and doing:

```sql
GRANT SELECT ON mytable TO PUBLIC;
GRANT SELECT, UPDATE, INSERT ON mytable TO admin;
GRANT SELECT (col1), UPDATE (col1) ON mytable TO miriam_rw;
```

For non-table objects there are other `\d` commands that can display their privileges.

If the “Access privileges” column is empty for a given object, it means the object has default privileges (that is, its privileges column is null). Default privileges always include all privileges for the owner, and can include some privileges for PUBLIC depending on the object type, as explained above. The first `GRANT` or `REVOKE` on an object will instantiate the default privileges (producing, for example, `{miriam=arwdDxt/miriam}`) and then modify them per the specified request. Similarly, entries are shown in “Column access privileges” only for columns with nondefault privileges. (Note: for this purpose, “default privileges” always means the built-in default privileges for the object's type. An object whose privileges have been affected by an `ALTER DEFAULT PRIVILEGES` command will always be shown with an explicit privilege entry that includes the effects of the `ALTER`.)

Notice that the owner's implicit grant options are not marked in the access privileges display. A `*` will appear only when grant options have been explicitly granted to someone.

### Examples

Grant insert privilege to all users on table `films`:

```sql
GRANT INSERT ON films TO PUBLIC;
```
GRANT INSERT ON films TO PUBLIC;

Grant all available privileges to user manuel on view kinds:

GRANT ALL PRIVILEGES ON kinds TO manuel;

Note that while the above will indeed grant all privileges if executed by a superuser or the owner of kinds, when executed by someone else it will only grant those permissions for which the someone else has grant options.

Grant membership in role admins to user joe:

GRANT admins TO joe;

Compatibility

According to the SQL standard, the PRIVILEGES key word in ALL PRIVILEGES is required. The SQL standard does not support setting the privileges on more than one object per command.

PostgreSQL allows an object owner to revoke their own ordinary privileges: for example, a table owner can make the table read-only to themselves by revoking their own INSERT, UPDATE, DELETE, and TRUNCATE privileges. This is not possible according to the SQL standard. The reason is that PostgreSQL treats the owner's privileges as having been granted by the owner to themselves; therefore they can revoke them too. In the SQL standard, the owner's privileges are granted by an assumed entity "_SYSTEM". Not being "_SYSTEM", the owner cannot revoke these rights.

According to the SQL standard, grant options can be granted to PUBLIC; PostgreSQL only supports granting grant options to roles.

The SQL standard provides for a USAGE privilege on other kinds of objects: character sets, collations, translations.

In the SQL standard, sequences only have a USAGE privilege, which controls the use of the NEXT VALUE FOR expression, which is equivalent to the function nextval in PostgreSQL. The sequence privileges SELECT and UPDATE are PostgreSQL extensions. The application of the sequence USAGE privilege to the currval function is also a PostgreSQL extension (as is the function itself).

Privileges on databases, tablespaces, schemas, and languages are PostgreSQL extensions.

See Also

REVOKE, ALTER DEFAULT PRIVILEGES
IMPORT FOREIGN SCHEMA

IMPORT FOREIGN SCHEMA — import table definitions from a foreign server

Synopsis

IMPORT FOREIGN SCHEMA remote_schema
    [ { LIMIT TO | EXCEPT } ( table_name [, ...] ) ]
FROM SERVER server_name
INTO local_schema
    [ OPTIONS ( option 'value' [, ... ] ) ]

Description

IMPORT FOREIGN SCHEMA creates foreign tables that represent tables existing on a foreign server.
The new foreign tables will be owned by the user issuing the command and are created with the correct
column definitions and options to match the remote tables.

By default, all tables and views existing in a particular schema on the foreign server are imported.
Optionally, the list of tables can be limited to a specified subset, or specific tables can be excluded.
The new foreign tables are all created in the target schema, which must already exist.

To use IMPORT FOREIGN SCHEMA, the user must have USAGE privilege on the foreign server, as
well as CREATE privilege on the target schema.

Parameters

remote_schema

The remote schema to import from. The specific meaning of a remote schema depends on the
foreign data wrapper in use.

LIMIT TO ( table_name [, ...] )

Import only foreign tables matching one of the given table names. Other tables existing in the
foreign schema will be ignored.

EXCEPT ( table_name [, ...] )

Exclude specified foreign tables from the import. All tables existing in the foreign schema will
be imported except the ones listed here.

server_name

The foreign server to import from.

local_schema

The schema in which the imported foreign tables will be created.

OPTIONS ( option 'value' [, ... ] )

Options to be used during the import. The allowed option names and values are specific to each
foreign data wrapper.
Examples

Import table definitions from a remote schema `foreign_films` on server `film_server`, creating the foreign tables in local schema `films`:

```
IMPORT FOREIGN SCHEMA foreign_films
    FROM SERVER film_server INTO films;
```

As above, but import only the two tables `actors` and `directors` (if they exist):

```
IMPORT FOREIGN SCHEMA foreign_films LIMIT TO (actors, directors)
    FROM SERVER film_server INTO films;
```

Compatibility

The `IMPORT FOREIGN SCHEMA` command conforms to the SQL standard, except that the `OPTIONS` clause is a PostgreSQL extension.

See Also

`CREATE FOREIGN TABLE`, `CREATE SERVER`
INSERT

INSERT — create new rows in a table

Synopsis

[ WITH [ RECURSIVE ] with_query [, ... ] ]
INSERT INTO table_name [ AS alias ] [ ( column_name [, ...] ) ]
[ OVERRIDING ( SYSTEM | USER) VALUE ]
( DEFAULT VALUES | VALUES ( { expression | DEFAULT } [, ...] )
[, ...] | query ]
[ ON CONFLICT [ conflict_target ] conflict_action ]
[ RETURNING * | output_expression [ [ AS ] output_name ]
[, ...] ]

where conflict_target can be one of:

( { index_column_name | ( index_expression
) } [ COLLATE collation ] [ opclass ] [, ...] )
WHERE index_predicate
ON CONSTRAINT constraint_name

and conflict_action is one of:

DO NOTHING
DO UPDATE SET { column_name = { expression | DEFAULT } |
( column_name [, ...] ) = [ ROW ]
( { expression | DEFAULT } [, ...] ) |
( column_name [, ...] ) = ( sub-SELECT )
} [, ...]
WHERE condition

Description

INSERT inserts new rows into a table. One can insert one or more rows specified by value expressions, or zero or more rows resulting from a query.

The target column names can be listed in any order. If no list of column names is given at all, the default is all the columns of the table in their declared order; or the first \( N \) column names, if there are only \( N \) columns supplied by the VALUES clause or query. The values supplied by the VALUES clause or query are associated with the explicit or implicit column list left-to-right.

Each column not present in the explicit or implicit column list will be filled with a default value, either its declared default value or null if there is none.

If the expression for any column is not of the correct data type, automatic type conversion will be attempted.

ON CONFLICT can be used to specify an alternative action to raising a unique constraint or exclusion constraint violation error. (See ON CONFLICT Clause below.)

The optional RETURNING clause causes INSERT to compute and return value(s) based on each row actually inserted (or updated, if an ON CONFLICT DO UPDATE clause was used). This is primarily useful for obtaining values that were supplied by defaults, such as a serial sequence number. However, any expression using the table's columns is allowed. The syntax of the RETURNING list is identical
to that of the output list of SELECT. Only rows that were successfully inserted or updated will be returned. For example, if a row was locked but not updated because an ON CONFLICT DO UPDATE ... WHERE clause condition was not satisfied, the row will not be returned.

You must have INSERT privilege on a table in order to insert into it. If ON CONFLICT DO UPDATE is present, UPDATE privilege on the table is also required.

If a column list is specified, you only need INSERT privilege on the listed columns. Similarly, when ON CONFLICT DO UPDATE is specified, you only need UPDATE privilege on the column(s) that are listed to be updated. However, ON CONFLICT DO UPDATE also requires SELECT privilege on any column whose values are read in the ON CONFLICT DO UPDATE expressions or condition.

Use of the RETURNING clause requires SELECT privilege on all columns mentioned in RETURNING. If you use the query clause to insert rows from a query, you of course need to have SELECT privilege on any table or column used in the query.

Parameters

Inserting

This section covers parameters that may be used when only inserting new rows. Parameters exclusively used with the ON CONFLICT clause are described separately.

with_query

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the INSERT query. See Section 7.8 and SELECT for details.

It is possible for the query (SELECT statement) to also contain a WITH clause. In such a case both sets of with_query can be referenced within the query, but the second one takes precedence since it is more closely nested.

table_name

The name (optionally schema-qualified) of an existing table.

alias

A substitute name for table_name. When an alias is provided, it completely hides the actual name of the table. This is particularly useful when ON CONFLICT DO UPDATE targets a table named excluded, since that will otherwise be taken as the name of the special table representing rows proposed for insertion.

column_name

The name of a column in the table named by table_name. The column name can be qualified with a subfield name or array subscript, if needed. (Inserting into only some fields of a composite column leaves the other fields null.) When referencing a column with ON CONFLICT DO UPDATE, do not include the table's name in the specification of a target column. For example, INSERT INTO table_name ... ON CONFLICT DO UPDATE SET table_name.col = 1 is invalid (this follows the general behavior for UPDATE).

OVERRIDING SYSTEM VALUE

Without this clause, it is an error to specify an explicit value (other than DEFAULT) for an identity column defined as GENERATED ALWAYS. This clause overrides that restriction.

OVERRIDING USER VALUE

If this clause is specified, then any values supplied for identity columns defined as GENERATED BY DEFAULT are ignored and the default sequence-generated values are applied.
This clause is useful for example when copying values between tables. Writing `INSERT INTO tbl2 OVERRIDING USER VALUE SELECT * FROM tbl1` will copy from `tbl1` all columns that are not identity columns in `tbl2` while values for the identity columns in `tbl2` will be generated by the sequences associated with `tbl2`.

**DEFAULT VALUES**

All columns will be filled with their default values. (An OVERRIDING clause is not permitted in this form.)

**expression**

An expression or value to assign to the corresponding column.

**DEFAULT**

The corresponding column will be filled with its default value.

**query**

A query (SELECT statement) that supplies the rows to be inserted. Refer to the SELECT statement for a description of the syntax.

**output_expression**

An expression to be computed and returned by the INSERT command after each row is inserted or updated. The expression can use any column names of the table named by `table_name`. Write `*` to return all columns of the inserted or updated row(s).

**output_name**

A name to use for a returned column.

**ON CONFLICT Clause**

The optional ON CONFLICT clause specifies an alternative action to raising a unique violation or exclusion constraint violation error. For each individual row proposed for insertion, either the insertion proceeds, or, if an arbiter constraint or index specified by `conflict_target` is violated, the alternative `conflict_action` is taken. ON CONFLICT DO NOTHING simply avoids inserting a row as its alternative action. ON CONFLICT DO UPDATE updates the existing row that conflicts with the row proposed for insertion as its alternative action.

`conflict_target` can perform unique index inference. When performing inference, it consists of one or more `index_column_name` columns and/or `index_expression` expressions, and an optional `index_predicate`. All `table_name` unique indexes that, without regard to order, contain exactly the `conflict_target`-specified columns/expressions are inferred (chosen) as arbiter indexes. If an `index_predicate` is specified, it must, as a further requirement for inference, satisfy arbiter indexes. Note that this means a non-partial unique index (a unique index without a predicate) will be inferred (and thus used by ON CONFLICT) if such an index satisfying every other criteria is available. If an attempt at inference is unsuccessful, an error is raised.

ON CONFLICT DO UPDATE guarantees an atomic INSERT or UPDATE outcome; provided there is no independent error, one of those two outcomes is guaranteed, even under high concurrency. This is also known as **UPSERT** — “UPDATE or INSERT”.

**conflict_target**

Specifies which conflicts ON CONFLICT takes the alternative action on by choosing arbiter indexes. Either performs unique index inference, or names a constraint explicitly. For ON CONFLICT DO NOTHING, it is optional to specify a `conflict_target`; when omitted, conflicts
with all usable constraints (and unique indexes) are handled. For ON CONFLICT DO UPDATE, a conflict_target must be provided.

conflict_action

conflict_action specifies an alternative ON CONFLICT action. It can be either DO NOTHING, or a DO UPDATE clause specifying the exact details of the UPDATE action to be performed in case of a conflict. The SET and WHERE clauses in ON CONFLICT DO UPDATE have access to the existing row using the table's name (or an alias), and to rows proposed for insertion using the special excluded table. SELECT privilege is required on any column in the target table where corresponding excluded columns are read.

Note that the effects of all per-row BEFORE INSERT triggers are reflected in excluded values, since those effects may have contributed to the row being excluded from insertion.

index_column_name

The name of a table_name column. Used to infer arbiter indexes. Follows CREATE INDEX format. SELECT privilege on index_column_name is required.

index_expression

Similar to index_column_name, but used to infer expressions on table_name columns appearing within index definitions (not simple columns). Follows CREATE INDEX format. SELECT privilege on any column appearing within index_expression is required.

collation

When specified, mandates that corresponding index_column_name or index_expression use a particular collation in order to be matched during inference. Typically this is omitted, as collations usually do not affect whether or not a constraint violation occurs. Follows CREATE INDEX format.

opclass

When specified, mandates that corresponding index_column_name or index_expression use particular operator class in order to be matched during inference. Typically this is omitted, as the equality semantics are often equivalent across a type's operator classes anyway, or because it's sufficient to trust that the defined unique indexes have the pertinent definition of equality. Follows CREATE INDEX format.

index_predicate

Used to allow inference of partial unique indexes. Any indexes that satisfy the predicate (which need not actually be partial indexes) can be inferred. Follows CREATE INDEX format. SELECT privilege on any column appearing within index_predicate is required.

constraint_name

Explicitly specifies an arbiter constraint by name, rather than inferring a constraint or index.

condition

An expression that returns a value of type boolean. Only rows for which this expression returns true will be updated, although all rows will be locked when the ON CONFLICT DO UPDATE action is taken. Note that condition is evaluated last, after a conflict has been identified as a candidate to update.

Note that exclusion constraints are not supported as arbiters with ON CONFLICT DO UPDATE. In all cases, only NOT DEFERRABLE constraints and unique indexes are supported as arbiters.
INSERT with an ON CONFLICT DO UPDATE clause is a “deterministic” statement. This means that the command will not be allowed to affect any single existing row more than once; a cardinality violation error will be raised when this situation arises. Rows proposed for insertion should not duplicate each other in terms of attributes constrained by an arbiter index or constraint.

Note that it is currently not supported for the ON CONFLICT DO UPDATE clause of an INSERT applied to a partitioned table to update the partition key of a conflicting row such that it requires the row be moved to a new partition.

### Tip
It is often preferable to use unique index inference rather than naming a constraint directly using ON CONFLICT ON CONSTRAINT constraint_name. Inference will continue to work correctly when the underlying index is replaced by another more or less equivalent index in an overlapping way, for example when using CREATE UNIQUE INDEX ... CONCURRENTLY before dropping the index being replaced.

### Outputs
On successful completion, an INSERT command returns a command tag of the form

```
INSERT oid count
```

The `count` is the number of rows inserted or updated. If `count` is exactly one, and the target table has OIDs, then `oid` is the OID assigned to the inserted row. The single row must have been inserted rather than updated. Otherwise `oid` is zero.

If the INSERT command contains a RETURNING clause, the result will be similar to that of a SELECT statement containing the columns and values defined in the RETURNING list, computed over the row(s) inserted or updated by the command.

### Notes
If the specified table is a partitioned table, each row is routed to the appropriate partition and inserted into it. If the specified table is a partition, an error will occur if one of the input rows violates the partition constraint.

### Examples
Insert a single row into table `films`:

```
INSERT INTO films VALUES
     ('UA502', 'Bananas', 105, '1971-07-13', 'Comedy', '82 minutes');
```

In this example, the `len` column is omitted and therefore it will have the default value:

```
INSERT INTO films (code, title, did, date_prod, kind)
    VALUES ('T_601', 'Yojimbo', 106, '1961-06-16', 'Drama');
```

This example uses the `DEFAULT` clause for the date columns rather than specifying a value:
INSERT INTO films VALUES ('UA502', 'Bananas', 105, DEFAULT, 'Comedy', '82 minutes');
INSERT INTO films (code, title, did, date_prod, kind) VALUES ('T_601', 'Yojimbo', 106, DEFAULT, 'Drama');

To insert a row consisting entirely of default values:

INSERT INTO films DEFAULT VALUES;

To insert multiple rows using the multirow VALUES syntax:

INSERT INTO films (code, title, did, date_prod, kind) VALUES ('B6717', 'Tampopo', 110, '1985-02-10', 'Comedy'), ('HG120', 'The Dinner Game', 140, DEFAULT, 'Comedy');

This example inserts some rows into table films from a table tmp_films with the same column layout as films:

INSERT INTO films SELECT * FROM tmp_films WHERE date_prod < '2004-05-07';

This example inserts into array columns:

-- Create an empty 3x3 gameboard for noughts-and-crosses
INSERT INTO tictactoe (game, board[1:3][1:3])
  VALUES (1, '{{"","",""},{"","",""},{"","",""}}');
-- The subscripts in the above example aren't really needed
INSERT INTO tictactoe (game, board)
  VALUES (2, '{{X,"",""},{"",O,""},{"",X,""}}');

Insert a single row into table distributors, returning the sequence number generated by the DEFAULT clause:

INSERT INTO distributors (did, dname) VALUES (DEFAULT, 'XYZ Widgets')
  RETURNING did;

Increment the sales count of the salesperson who manages the account for Acme Corporation, and record the whole updated row along with current time in a log table:

WITH upd AS (
    UPDATE employees SET sales_count = sales_count + 1 WHERE id =
    (SELECT sales_person FROM accounts WHERE name = 'Acme Corporation')
    RETURNING *
)
INSERT INTO employees_log SELECT *, current_timestamp FROM upd;

Insert or update new distributors as appropriate. Assumes a unique index has been defined that constrains values appearing in the did column. Note that the special excluded table is used to reference values originally proposed for insertion:

INSERT INTO distributors (did, dname)
VALUES (5, 'Gizmo Transglobal'), (6, 'Associated Computing, Inc')
    ON CONFLICT (did) DO UPDATE SET dname = EXCLUDED.dname;

Insert a distributor, or do nothing for rows proposed for insertion when an existing, excluded row (a row with a matching constrained column or columns after before row insert triggers fire) exists. Example assumes a unique index has been defined that constrains values appearing in the did column:

    INSERT INTO distributors (did, dname) VALUES (7, 'Redline GmbH')
    ON CONFLICT (did) DO NOTHING;

Insert or update new distributors as appropriate. Example assumes a unique index has been defined that constrains values appearing in the did column. WHERE clause is used to limit the rows actually updated (any existing row not updated will still be locked, though):

    -- Don't update existing distributors based in a certain ZIP code
    INSERT INTO distributors AS d (did, dname) VALUES (8, 'Anvil Distribution')
    ON CONFLICT (did) DO UPDATE
        SET dname = EXCLUDED.dname || ' (formerly ' || d.dname || ')' WHERE d.zipcode <> '21201';

    -- Name a constraint directly in the statement (uses associated -- index to arbitrate taking the DO NOTHING action)
    INSERT INTO distributors (did, dname) VALUES (9, 'Antwerp Design')
    ON CONFLICT ON CONSTRAINT distributors_pkey DO NOTHING;

    -- Insert new distributor if possible; otherwise DO NOTHING. Example assumes a unique index has been defined that constrains values appearing in the did column on a subset of rows where the is_active Boolean column evaluates to true:

    -- This statement could infer a partial unique index on "did"
    -- with a predicate of "WHERE is_active", but it could also
    -- just use a regular unique constraint on "did"
    INSERT INTO distributors (did, dname) VALUES (10, 'Conrad International')
    ON CONFLICT (did) WHERE is_active DO NOTHING;

Compatibility

INSERT conforms to the SQL standard, except that the RETURNING clause is a PostgreSQL extension, as is the ability to use WITH with INSERT, and the ability to specify an alternative action with ON CONFLICT. Also, the case in which a column name list is omitted, but not all the columns are filled from the VALUES clause or query, is disallowed by the standard.

The SQL standard specifies that OVERRIDING SYSTEM VALUE can only be specified if an identity column that is generated always exists. PostgreSQL allows the clause in any case and ignores it if it is not applicable.

Possible limitations of the query clause are documented under SELECT.
LISTEN

LISTEN — listen for a notification

Synopsis

LISTEN channel

Description

LISTEN registers the current session as a listener on the notification channel named channel. If the current session is already registered as a listener for this notification channel, nothing is done.

Whenever the command NOTIFY channel is invoked, either by this session or another one connected to the same database, all the sessions currently listening on that notification channel are notified, and each will in turn notify its connected client application.

A session can be unregistered for a given notification channel with the UNLISTEN command. A session's listen registrations are automatically cleared when the session ends.

The method a client application must use to detect notification events depends on which PostgreSQL application programming interface it uses. With the libpq library, the application issues LISTEN as an ordinary SQL command, and then must periodically call the function PQnotifies to find out whether any notification events have been received. Other interfaces such as libpq provide higher-level methods for handling notify events; indeed, with libpq the application programmer should not even issue LISTEN or UNLISTEN directly. See the documentation for the interface you are using for more details.

NOTIFY contains a more extensive discussion of the use of LISTEN and NOTIFY.

Parameters

channel

Name of a notification channel (any identifier).

Notes

LISTEN takes effect at transaction commit. If LISTEN or UNLISTEN is executed within a transaction that later rolls back, the set of notification channels being listened to is unchanged.

A transaction that has executed LISTEN cannot be prepared for two-phase commit.

Examples

Configure and execute a listen/notify sequence from psql:

LISTEN virtual;
NOTIFY virtual;
Asynchronous notification "virtual" received from server process with PID 8448.
Compatibility

There is no `LISTEN` statement in the SQL standard.

See Also

`NOTIFY`, `UNLISTEN`
LOAD

LOAD — load a shared library file

Synopsis

LOAD 'filename'

Description

This command loads a shared library file into the PostgreSQL server’s address space. If the file has been loaded already, the command does nothing. Shared library files that contain C functions are automatically loaded whenever one of their functions is called. Therefore, an explicit LOAD is usually only needed to load a library that modifies the server’s behavior through “hooks” rather than providing a set of functions.

The library file name is typically given as just a bare file name, which is sought in the server’s library search path (set by dynamic_library_path). Alternatively it can be given as a full path name. In either case the platform’s standard shared library file name extension may be omitted. See Section 38.10.1 for more information on this topic.

Non-superusers can only apply LOAD to library files located in $libdir/plugins/ — the specified filename must begin with exactly that string. (It is the database administrator’s responsibility to ensure that only “safe” libraries are installed there.)

Compatibility

LOAD is a PostgreSQL extension.

See Also

CREATE FUNCTION
LOCK

LOCK — lock a table

Synopsis

LOCK [ TABLE ] [ ONLY ] name [ * ] [, ...] [ IN lockmode MODE ]
[ NOWAIT ]

where lockmode is one of:

ACCESS SHARE | ROW SHARE | ROW EXCLUSIVE | SHARE UPDATE
EXCLUSIVE
| SHARE | SHARE ROW EXCLUSIVE | EXCLUSIVE | ACCESS EXCLUSIVE

Description

LOCK TABLE obtains a table-level lock, waiting if necessary for any conflicting locks to be released. If NOWAIT is specified, LOCK TABLE does not wait to acquire the desired lock: if it cannot be acquired immediately, the command is aborted and an error is emitted. Once obtained, the lock is held for the remainder of the current transaction. (There is no UNLOCK TABLE command; locks are always released at transaction end.)

When a view is locked, all relations appearing in the view definition query are also locked recursively with the same lock mode.

When acquiring locks automatically for commands that reference tables, PostgreSQL always uses the least restrictive lock mode possible. LOCK TABLE provides for cases when you might need more restrictive locking. For example, suppose an application runs a transaction at the READ COMMITTED isolation level and needs to ensure that data in a table remains stable for the duration of the transaction. To achieve this you could obtain SHARE lock mode over the table before querying. This will prevent concurrent data changes and ensure subsequent reads of the table see a stable view of committed data, because SHARE lock mode conflicts with the ROW EXCLUSIVE lock acquired by writers, and your LOCK TABLE name IN SHARE MODE statement will wait until any concurrent holders of ROW EXCLUSIVE mode locks commit or roll back. Thus, once you obtain the lock, there are no uncommitted writes outstanding; furthermore none can begin until you release the lock.

To achieve a similar effect when running a transaction at the REPEATABLE READ or SERIALIZABLE isolation level, you have to execute the LOCK TABLE statement before executing any SELECT or data modification statement. A REPEATABLE READ or SERIALIZABLE transaction's view of data will be frozen when its first SELECT or data modification statement begins. A LOCK TABLE later in the transaction will still prevent concurrent writes — but it won't ensure that what the transaction reads corresponds to the latest committed values.

If a transaction of this sort is going to change the data in the table, then it should use SHARE ROW EXCLUSIVE lock mode instead of SHARE mode. This ensures that only one transaction of this type runs at a time. Without this, a deadlock is possible: two transactions might both acquire SHARE mode, and then be unable to also acquire ROW EXCLUSIVE mode to actually perform their updates. (Note that a transaction's own locks never conflict, so a transaction can acquire ROW EXCLUSIVE mode when it holds SHARE mode — but not if anyone else holds SHARE mode.) To avoid deadlocks, make sure all transactions acquire locks on the same objects in the same order, and if multiple lock modes are involved for a single object, then transactions should always acquire the most restrictive mode first.

More information about the lock modes and locking strategies can be found in Section 13.3.
Parameters

name

The name (optionally schema-qualified) of an existing table to lock. If ONLY is specified before the table name, only that table is locked. If ONLY is not specified, the table and all its descendant tables (if any) are locked. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

The command LOCK TABLE a, b; is equivalent to LOCK TABLE a; LOCK TABLE b;.

The tables are locked one-by-one in the order specified in the LOCK TABLE command.

lockmode

The lock mode specifies which locks this lock conflicts with. Lock modes are described in Section 13.3.

If no lock mode is specified, then ACCESS EXCLUSIVE, the most restrictive mode, is used.

NOWAIT

Specifies that LOCK TABLE should not wait for any conflicting locks to be released: if the specified lock(s) cannot be acquired immediately without waiting, the transaction is aborted.

Notes

LOCK TABLE ... IN ACCESS SHARE MODE requires SELECT privileges on the target table. LOCK TABLE ... IN ROW EXCLUSIVE MODE requires INSERT, UPDATE, DELETE, or TRUNCATE privileges on the target table. All other forms of LOCK require table-level UPDATE, DELETE, or TRUNCATE privileges.

The user performing the lock on the view must have the corresponding privilege on the view. In addition the view’s owner must have the relevant privileges on the underlying base relations, but the user performing the lock does not need any permissions on the underlying base relations.

LOCK TABLE is useless outside a transaction block: the lock would remain held only to the completion of the statement. Therefore PostgreSQL reports an error if LOCK is used outside a transaction block. Use BEGIN and COMMIT (or ROLLBACK) to define a transaction block.

LOCK TABLE only deals with table-level locks, and so the mode names involving ROW are all misnomers. These mode names should generally be read as indicating the intention of the user to acquire row-level locks within the locked table. Also, ROW EXCLUSIVE mode is a shareable table lock. Keep in mind that all the lock modes have identical semantics so far as LOCK TABLE is concerned, differing only in the rules about which modes conflict with which. For information on how to acquire an actual row-level lock, see Section 13.3.2 and the The Locking Clause in the SELECT reference documentation.

Examples

Obtain a SHARE lock on a primary key table when going to perform inserts into a foreign key table:

BEGIN WORK;
LOCK TABLE films IN SHARE MODE;
SELECT id FROM films
    WHERE name = 'Star Wars: Episode I - The Phantom Menace';
-- Do ROLLBACK if record was not returned
INSERT INTO films_user_comments VALUES
    (_id_, 'GREAT! I was waiting for it for so long!');
LOCK

COMMIT WORK;

Take a SHARE ROW EXCLUSIVE lock on a primary key table when going to perform a delete operation:

BEGIN WORK;
LOCK TABLE films IN SHARE ROW EXCLUSIVE MODE;
DELETE FROM films_user_comments WHERE id IN
   (SELECT id FROM films WHERE rating < 5);
DELETE FROM films WHERE rating < 5;
COMMIT WORK;

Compatibility

There is no LOCK TABLE in the SQL standard, which instead uses SET TRANSACTION to specify concurrency levels on transactions. PostgreSQL supports that too; see SET TRANSACTION for details.

Except for ACCESS SHARE, ACCESS EXCLUSIVE, and SHARE UPDATE EXCLUSIVE lock modes, the PostgreSQL lock modes and the LOCK TABLE syntax are compatible with those present in Oracle.
**MOVE**

MOVE — position a cursor

**Synopsis**

MOVE [ direction [ FROM | IN ] ] cursor_name

where direction can be empty or one of:

- NEXT
- PRIOR
- FIRST
- LAST
- ABSOLUTE count
- RELATIVE count
- count
- ALL
- FORWARD
- FORWARD count
- FORWARD ALL
- BACKWARD
- BACKWARD count
- BACKWARD ALL

**Description**

MOVE repositions a cursor without retrieving any data. MOVE works exactly like the FETCH command, except it only positions the cursor and does not return rows.

The parameters for the MOVE command are identical to those of the FETCH command; refer to FETCH for details on syntax and usage.

**Outputs**

On successful completion, a MOVE command returns a command tag of the form

MOVE count

The count is the number of rows that a FETCH command with the same parameters would have returned (possibly zero).

**Examples**

BEGIN WORK;
DECLARE liahona CURSOR FOR SELECT * FROM films;

-- Skip the first 5 rows:
MOVE FORWARD 5 IN liahona;
MOVE 5

-- Fetch the 6th row from the cursor liahona:
MOVE

FETCH 1 FROM liahoma;

<table>
<thead>
<tr>
<th>code</th>
<th>title</th>
<th>did</th>
<th>date_prod</th>
<th>kind</th>
<th>len</th>
</tr>
</thead>
<tbody>
<tr>
<td>P_303</td>
<td>48 Hrs</td>
<td>103</td>
<td>1982-10-22</td>
<td>Action</td>
<td>01:37</td>
</tr>
</tbody>
</table>

(1 row)

-- Close the cursor liahoma and end the transaction:
CLOSE liahoma;
COMMIT WORK;

Compatibility

There is no MOVE statement in the SQL standard.

See Also

CLOSE, DECLARE, FETCH
NOTIFY

NOTIFY — generate a notification

Synopsis

NOTIFY channel [ , payload ]

Description

The NOTIFY command sends a notification event together with an optional “payload” string to each client application that has previously executed LISTEN channel for the specified channel name in the current database. Notifications are visible to all users.

NOTIFY provides a simple interprocess communication mechanism for a collection of processes accessing the same PostgreSQL database. A payload string can be sent along with the notification, and higher-level mechanisms for passing structured data can be built by using tables in the database to pass additional data from notifier to listener(s).

The information passed to the client for a notification event includes the notification channel name, the notifying session's server process PID, and the payload string, which is an empty string if it has not been specified.

It is up to the database designer to define the channel names that will be used in a given database and what each one means. Commonly, the channel name is the same as the name of some table in the database, and the notify event essentially means, “I changed this table, take a look at it to see what's new”. But no such association is enforced by the NOTIFY and LISTEN commands. For example, a database designer could use several different channel names to signal different sorts of changes to a single table. Alternatively, the payload string could be used to differentiate various cases.

When NOTIFY is used to signal the occurrence of changes to a particular table, a useful programming technique is to put the NOTIFY in a statement trigger that is triggered by table updates. In this way, notification happens automatically when the table is changed, and the application programmer cannot accidentally forget to do it.

NOTIFY interacts with SQL transactions in some important ways. Firstly, if a NOTIFY is executed inside a transaction, the notify events are not delivered until and unless the transaction is committed. This is appropriate, since if the transaction is aborted, all the commands within it have had no effect, including NOTIFY. But it can be disconcerting if one is expecting the notification events to be delivered immediately. Secondly, if a listening session receives a notification signal while it is within a transaction, the notification event will not be delivered to its connected client until just after the transaction is completed (either committed or aborted). Again, the reasoning is that if a notification were delivered within a transaction that was later aborted, one would want the notification to be undone somehow — but the server cannot “take back” a notification once it has sent it to the client. So notification events are only delivered between transactions. The upshot of this is that applications using NOTIFY for real-time signaling should try to keep their transactions short.

If the same channel name is signaled multiple times from the same transaction with identical payload strings, the database server can decide to deliver a single notification only. On the other hand, notifications with distinct payload strings will always be delivered as distinct notifications. Similarly, notifications from different transactions will never get folded into one notification. Except for dropping later instances of duplicate notifications, NOTIFY guarantees that notifications from the same transaction get delivered in the order they were sent. It is also guaranteed that messages from different transactions are delivered in the order in which the transactions committed.
It is common for a client that executes NOTIFY to be listening on the same notification channel itself. In that case it will get back a notification event, just like all the other listening sessions. Depending on the application logic, this could result in useless work, for example, reading a database table to find the same updates that that session just wrote out. It is possible to avoid such extra work by noticing whether the notifying session’s server process PID (supplied in the notification event message) is the same as one’s own session’s PID (available from libpq). When they are the same, the notification event is one’s own work bouncing back, and can be ignored.

**Parameters**

*channel*

Name of the notification channel to be signaled (any identifier).

*payload*

The “payload” string to be communicated along with the notification. This must be specified as a simple string literal. In the default configuration it must be shorter than 8000 bytes. (If binary data or large amounts of information need to be communicated, it's best to put it in a database table and send the key of the record.)

**Notes**

There is a queue that holds notifications that have been sent but not yet processed by all listening sessions. If this queue becomes full, transactions calling NOTIFY will fail at commit. The queue is quite large (8GB in a standard installation) and should be sufficiently sized for almost every use case. However, no cleanup can take place if a session executes LISTEN and then enters a transaction for a very long time. Once the queue is half full you will see warnings in the log file pointing you to the session that is preventing cleanup. In this case you should make sure that this session ends its current transaction so that cleanup can proceed.

The function `pg_notification_queue_usage` returns the fraction of the queue that is currently occupied by pending notifications. See Section 9.25 for more information.

A transaction that has executed NOTIFY cannot be prepared for two-phase commit.

**pg_notify**

To send a notification you can also use the function `pg_notify(text, text)`. The function takes the channel name as the first argument and the payload as the second. The function is much easier to use than the NOTIFY command if you need to work with non-constant channel names and payloads.

**Examples**

Configure and execute a listen/notify sequence from psql:

```sql
LISTEN virtual;
NOTIFY virtual;
Asynchronous notification "virtual" received from server process with PID 8448.
NOTIFY virtual, 'This is the payload';
Asynchronous notification "virtual" with payload "This is the payload" received from server process with PID 8448.

LISTEN foo;
SELECT pg_notify('fo' || 'o', 'pay' || 'load');
```
Asynchronous notification "foo" with payload "payload" received from server process with PID 14728.

Compatibility

There is no NOTIFY statement in the SQL standard.

See Also

LISTEN, UNLISTEN
PREPARE

PREPARE — prepare a statement for execution

Synopsis

PREPARE name [ ( data_type [, ...] ) ] AS statement

Description

PREPARE creates a prepared statement. A prepared statement is a server-side object that can be used to optimize performance. When the PREPARE statement is executed, the specified statement is parsed, analyzed, and rewritten. When an EXECUTE command is subsequently issued, the prepared statement is planned and executed. This division of labor avoids repetitive parse analysis work, while allowing the execution plan to depend on the specific parameter values supplied.

Prepared statements can take parameters: values that are substituted into the statement when it is executed. When creating the prepared statement, refer to parameters by position, using $1, $2, etc. A corresponding list of parameter data types can optionally be specified. When a parameter's data type is not specified or is declared as unknown, the type is inferred from the context in which the parameter is first referenced (if possible). When executing the statement, specify the actual values for these parameters in the EXECUTE statement. Refer to EXECUTE for more information about that.

Prepared statements only last for the duration of the current database session. When the session ends, the prepared statement is forgotten, so it must be recreated before being used again. This also means that a single prepared statement cannot be used by multiple simultaneous database clients; however, each client can create their own prepared statement to use. Prepared statements can be manually cleaned up using the DEALLOCATE command.

Prepared statements potentially have the largest performance advantage when a single session is being used to execute a large number of similar statements. The performance difference will be particularly significant if the statements are complex to plan or rewrite, e.g. if the query involves a join of many tables or requires the application of several rules. If the statement is relatively simple to plan and rewrite but relatively expensive to execute, the performance advantage of prepared statements will be less noticeable.

Parameters

name

An arbitrary name given to this particular prepared statement. It must be unique within a single session and is subsequently used to execute or deallocate a previously prepared statement.

data_type

The data type of a parameter to the prepared statement. If the data type of a particular parameter is unspecified or is specified as unknown, it will be inferred from the context in which the parameter is first referenced. To refer to the parameters in the prepared statement itself, use $1, $2, etc.

statement

Any SELECT, INSERT, UPDATE, DELETE, or VALUES statement.

Notes

Prepared statements can use generic plans rather than re-planning with each set of supplied EXECUTE values. This occurs immediately for prepared statements with no parameters; otherwise it occurs only
after five or more executions produce plans whose estimated cost average (including planning overhead) is more expensive than the generic plan cost estimate. Once a generic plan is chosen, it is used for the remaining lifetime of the prepared statement. Using EXECUTE values which are rare in columns with many duplicates can generate custom plans that are so much cheaper than the generic plan, even after adding planning overhead, that the generic plan might never be used.

A generic plan assumes that each value supplied to EXECUTE is one of the column’s distinct values and that column values are uniformly distributed. For example, if statistics record three distinct column values, a generic plan assumes a column equality comparison will match 33% of processed rows. Column statistics also allow generic plans to accurately compute the selectivity of unique columns. Comparisons on non-uniformly-distributed columns and specification of non-existent values affects the average plan cost, and hence if and when a generic plan is chosen.

To examine the query plan PostgreSQL is using for a prepared statement, use EXPLAIN, e.g. EXPLAIN EXECUTE. If a generic plan is in use, it will contain parameter symbols $n$, while a custom plan will have the supplied parameter values substituted into it. The row estimates in the generic plan reflect the selectivity computed for the parameters.

For more information on query planning and the statistics collected by PostgreSQL for that purpose, see the ANALYZE documentation.

Although the main point of a prepared statement is to avoid repeated parse analysis and planning of the statement, PostgreSQL will force re-analysis and re-planning of the statement before using it whenever database objects used in the statement have undergone definitional (DDL) changes since the previous use of the prepared statement. Also, if the value of search_path changes from one use to the next, the statement will be re-parsed using the new search_path. (This latter behavior is new as of PostgreSQL 9.3.) These rules make use of a prepared statement semantically almost equivalent to re-submitting the same query text over and over, but with a performance benefit if no object definitions are changed, especially if the best plan remains the same across uses. An example of a case where the semantic equivalence is not perfect is that if the statement refers to a table by an unqualified name, and then a new table of the same name is created in a schema appearing earlier in the search_path, no automatic re-parse will occur since no object used in the statement changed. However, if some other change forces a re-parse, the new table will be referenced in subsequent uses.

You can see all prepared statements available in the session by querying the pg_prepared_statements system view.

**Examples**

Create a prepared statement for an INSERT statement, and then execute it:

```
PREPARE fooplan (int, text, bool, numeric) AS
    INSERT INTO foo VALUES($1, $2, $3, $4);
EXECUTE fooplan(1, 'Hunter Valley', 't', 200.00);
```

Create a prepared statement for a SELECT statement, and then execute it:

```
PREPARE usrrptplan (int) AS
    SELECT * FROM users u, logs l WHERE u.usrid=$1 AND u.usrid=l.usrid
    AND l.date = $2;
EXECUTE usrrptplan(1, current_date);
```

Note that the data type of the second parameter is not specified, so it is inferred from the context in which $2 is used.
PREPARE

Compatibility

The SQL standard includes a PREPARE statement, but it is only for use in embedded SQL. This version of the PREPARE statement also uses a somewhat different syntax.

See Also

DEALLOCATE, EXECUTE
PREPARE TRANSACTION

PREPARE TRANSACTION — prepare the current transaction for two-phase commit

Synopsis

PREPARE TRANSACTION transaction_id

Description

PREPARE TRANSACTION prepares the current transaction for two-phase commit. After this command, the transaction is no longer associated with the current session; instead, its state is fully stored on disk, and there is a very high probability that it can be committed successfully, even if a database crash occurs before the commit is requested.

Once prepared, a transaction can later be committed or rolled back with COMMIT PREPARED or ROLLBACK PREPARED, respectively. Those commands can be issued from any session, not only the one that executed the original transaction.

From the point of view of the issuing session, PREPARE TRANSACTION is not unlike a ROLLBACK command: after executing it, there is no active current transaction, and the effects of the prepared transaction are no longer visible. (The effects will become visible again if the transaction is committed.)

If the PREPARE TRANSACTION command fails for any reason, it becomes a ROLLBACK: the current transaction is canceled.

Parameters

transaction_id

An arbitrary identifier that later identifies this transaction for COMMIT PREPARED or ROLLBACK PREPARED. The identifier must be written as a string literal, and must be less than 200 bytes long. It must not be the same as the identifier used for any currently prepared transaction.

Notes

PREPARE TRANSACTION is not intended for use in applications or interactive sessions. Its purpose is to allow an external transaction manager to perform atomic global transactions across multiple databases or other transactional resources. Unless you're writing a transaction manager, you probably shouldn't be using PREPARE TRANSACTION.

This command must be used inside a transaction block. Use BEGIN to start one.

It is not currently allowed to PREPARE a transaction that has executed any operations involving temporary tables, created any cursors WITH HOLD, or executed LISTEN, UNLISTEN, or NOTIFY. Those features are too tightly tied to the current session to be useful in a transaction to be prepared.

If the transaction modified any run-time parameters with SET (without the LOCAL option), those effects persist after PREPARE TRANSACTION, and will not be affected by any later COMMIT PREPARED or ROLLBACK PREPARED. Thus, in this one respect PREPARE TRANSACTION acts more like COMMIT than ROLLBACK.

All currently available prepared transactions are listed in the pg_prepared_xacts system view.
Caution

It is unwise to leave transactions in the prepared state for a long time. This will interfere with the ability of VACUUM to reclaim storage, and in extreme cases could cause the database to shut down to prevent transaction ID wraparound (see Section 24.1.5). Keep in mind also that the transaction continues to hold whatever locks it held. The intended usage of the feature is that a prepared transaction will normally be committed or rolled back as soon as an external transaction manager has verified that other databases are also prepared to commit.

If you have not set up an external transaction manager to track prepared transactions and ensure they get closed out promptly, it is best to keep the prepared-transaction feature disabled by setting max_prepared_transactions to zero. This will prevent accidental creation of prepared transactions that might then be forgotten and eventually cause problems.

Examples

Prepare the current transaction for two-phase commit, using foobar as the transaction identifier:

PREPARE TRANSACTION 'foobar';

Compatibility

PREPARE TRANSACTION is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.

See Also

COMMIT PREPARED, ROLLBACK PREPARED
REASSIGN OWNED

REASSIGN OWNED — change the ownership of database objects owned by a database role

Synopsis

REASSIGN OWNED BY { old_role | CURRENT_USER | SESSION_USER }
[,...]
TO { new_role | CURRENT_USER | SESSION_USER }

Description

REASSIGN OWNED instructs the system to change the ownership of database objects owned by any of the old_roles to new_role.

Parameters

old_role

The name of a role. The ownership of all the objects within the current database, and of all shared objects (databases, tablespaces), owned by this role will be reassigned to new_role.

new_role

The name of the role that will be made the new owner of the affected objects.

Notes

REASSIGN OWNED is often used to prepare for the removal of one or more roles. Because REASSIGN OWNED does not affect objects within other databases, it is usually necessary to execute this command in each database that contains objects owned by a role that is to be removed.

REASSIGN OWNED requires membership on both the source role(s) and the target role.

The DROP OWNED command is an alternative that simply drops all the database objects owned by one or more roles.

The REASSIGN OWNED command does not affect any privileges granted to the old_roles for objects that are not owned by them. Use DROP OWNED to revoke such privileges.

See Section 21.4 for more discussion.

Compatibility

The REASSIGN OWNED command is a PostgreSQL extension.

See Also

DROP OWNED, DROP ROLE, ALTER DATABASE
REFRESH MATERIALIZED VIEW

REFRESH MATERIALIZED VIEW — replace the contents of a materialized view

Synopsis

REFRESH MATERIALIZED VIEW [ CONCURRENTLY ] name
[ WITH [ NO ] DATA ]

Description

REFRESH MATERIALIZED VIEW completely replaces the contents of a materialized view. To execute this command you must be the owner of the materialized view. The old contents are discarded. If WITH DATA is specified (or defaults) the backing query is executed to provide the new data, and the materialized view is left in a scannable state. If WITH NO DATA is specified no new data is generated and the materialized view is left in an unscannable state.

CONCURRENTLY and WITH NO DATA may not be specified together.

Parameters

CONCURRENTLY

Refresh the materialized view without locking out concurrent selects on the materialized view. Without this option a refresh which affects a lot of rows will tend to use fewer resources and complete more quickly, but could block other connections which are trying to read from the materialized view. This option may be faster in cases where a small number of rows are affected.

This option is only allowed if there is at least one UNIQUE index on the materialized view which uses only column names and includes all rows; that is, it must not index on any expressions nor include a WHERE clause.

This option may not be used when the materialized view is not already populated.

Even with this option only one REFRESH at a time may run against any one materialized view.

name

The name (optionally schema-qualified) of the materialized view to refresh.

Notes

While the default index for future CLUSTER operations is retained, REFRESH MATERIALIZED VIEW does not order the generated rows based on this property. If you want the data to be ordered upon generation, you must use an ORDER BY clause in the backing query.

Examples

This command will replace the contents of the materialized view called order_summary using the query from the materialized view’s definition, and leave it in a scannable state:

REFRESH MATERIALIZED VIEW order_summary;
This command will free storage associated with the materialized view `annual_statistics_basis` and leave it in an unscannable state:

```
REFRESH MATERIALIZED VIEW annual_statistics_basis WITH NO DATA;
```

**Compatibility**

`REFRESH MATERIALIZED VIEW` is a PostgreSQL extension.

**See Also**

CREATE MATERIALIZED VIEW, ALTER MATERIALIZED VIEW, DROP MATERIALIZED VIEW
REINDEX

REINDEX — rebuild indexes

Synopsis

REINDEX [ ( VERBOSE ) ] { INDEX | TABLE | SCHEMA | DATABASE | SYSTEM } name

Description

REINDEX rebuilds an index using the data stored in the index's table, replacing the old copy of the index. There are several scenarios in which to use REINDEX:

- An index has become corrupted, and no longer contains valid data. Although in theory this should never happen, in practice indexes can become corrupted due to software bugs or hardware failures. REINDEX provides a recovery method.

- An index has become "bloated", that is it contains many empty or nearly-empty pages. This can occur with B-tree indexes in PostgreSQL under certain uncommon access patterns. REINDEX provides a way to reduce the space consumption of the index by writing a new version of the index without the dead pages. See Section 24.2 for more information.

- You have altered a storage parameter (such as fillfactor) for an index, and wish to ensure that the change has taken full effect.

- An index build with the CONCURRENTLY option failed, leaving an “invalid” index. Such indexes are useless but it can be convenient to use REINDEX to rebuild them. Note that REINDEX will not perform a concurrent build. To build the index without interfering with production you should drop the index and reissue the CREATE INDEX CONCURRENTLY command.

Parameters

INDEX

Recreate the specified index.

TABLE

Recreate all indexes of the specified table. If the table has a secondary “TOAST” table, that is reindexed as well.

SCHEMA

Recreate all indexes of the specified schema. If a table of this schema has a secondary “TOAST” table, that is reindexed as well. Indexes on shared system catalogs are also processed. This form of REINDEX cannot be executed inside a transaction block.

DATABASE

Recreate all indexes within the current database. Indexes on shared system catalogs are also processed. This form of REINDEX cannot be executed inside a transaction block.

SYSTEM

Recreate all indexes on system catalogs within the current database. Indexes on shared system catalogs are included. Indexes on user tables are not processed. This form of REINDEX cannot be executed inside a transaction block.
name

The name of the specific index, table, or database to be reindexed. Index and table names can be
schema-qualified. Presently, \texttt{REINDEX DATABASE} and \texttt{REINDEX SYSTEM} can only reindex
the current database, so their parameter must match the current database's name.

VERBOSE

Prints a progress report as each index is reindexed.

Notes

If you suspect corruption of an index on a user table, you can simply rebuild that index, or all indexes
on the table, using \texttt{REINDEX INDEX} or \texttt{REINDEX TABLE}.

Things are more difficult if you need to recover from corruption of an index on a system table. In
this case it's important for the system to not have used any of the suspect indexes itself. (Indeed, in
this sort of scenario you might find that server processes are crashing immediately at start-up, due to
reliance on the corrupted indexes.) To recover safely, the server must be started with the \texttt{-P} option,
which prevents it from using indexes for system catalog lookups.

One way to do this is to shut down the server and start a single-user PostgreSQL server with the \texttt{-P}
option included on its command line. Then, \texttt{REINDEX DATABASE}, \texttt{REINDEX SYSTEM}, \texttt{REINDEX}
TABLE, or \texttt{REINDEX INDEX} can be issued, depending on how much you want to reconstruct. If in
doubt, use \texttt{REINDEX SYSTEM} to select reconstruction of all system indexes in the database. Then
quit the single-user server session and restart the regular server. See the postgres reference page for
more information about how to interact with the single-user server interface.

Alternatively, a regular server session can be started with \texttt{-P} included in its command line options.
The method for doing this varies across clients, but in all libpq-based clients, it is possible to set the
\texttt{POPTIONS} environment variable to \texttt{-P} before starting the client. Note that while this method does
not require locking out other clients, it might still be wise to prevent other users from connecting to
the damaged database until repairs have been completed.

\texttt{REINDEX} is similar to a drop and recreate of the index in that the index contents are rebuilt from
scratch. However, the locking considerations are rather different. \texttt{REINDEX} locks out writes but not
reads of the index's parent table. It also takes an exclusive lock on the specific index being processed,
which will block reads that attempt to use that index. In contrast, \texttt{DROP INDEX} momentarily takes an
exclusive lock on the parent table, blocking both writes and reads. The subsequent \texttt{CREATE INDEX}
locks out writes but not reads; since the index is not there, no read will attempt to use it, meaning that
there will be no blocking but reads might be forced into expensive sequential scans.

Reindexing a single index or table requires being the owner of that index or table. Reindexing a schema
or database requires being the owner of that schema or database. Note that is therefore sometimes
possible for non-superusers to rebuild indexes of tables owned by other users. However, as a special
exception, when \texttt{REINDEX DATABASE}, \texttt{REINDEX SCHEMA} or \texttt{REINDEX SYSTEM} is issued by
a non-superuser, indexes on shared catalogs will be skipped unless the user owns the catalog (which
typically won't be the case). Of course, superusers can always reindex anything.

Reindexing partitioned tables or partitioned indexes is not supported. Each individual partition can be
reindexed separately instead.

Examples

Rebuild a single index:

\begin{verbatim}
REINDEX INDEX my_index;
\end{verbatim}

Rebuild all the indexes on the table \texttt{my_table}:
REINDEX TABLE my_table;

Rebuild all indexes in a particular database, without trusting the system indexes to be valid already:

$ export PGOPTIONS="-P"
$ psql broken_db
...
broken_db=> REINDEX DATABASE broken_db;
broken_db=> \q

Compatibility

There is no REINDEX command in the SQL standard.
RELEASE SAVEPOINT

RELEASE SAVEPOINT — destroy a previously defined savepoint

Synopsis

RELEASE [ SAVEPOINT ] savepoint_name

Description

RELEASE SAVEPOINT destroys a savepoint previously defined in the current transaction.

Destroying a savepoint makes it unavailable as a rollback point, but it has no other user visible behavior. It does not undo the effects of commands executed after the savepoint was established. (To do that, see ROLLBACK TO SAVEPOINT.) Destroying a savepoint when it is no longer needed allows the system to reclaim some resources earlier than transaction end.

RELEASE SAVEPOINT also destroys all savepoints that were established after the named savepoint was established.

Parameters

savepoint_name

The name of the savepoint to destroy.

Notes

Specifying a savepoint name that was not previously defined is an error.

It is not possible to release a savepoint when the transaction is in an aborted state.

If multiple savepoints have the same name, only the one that was most recently defined is released.

Examples

To establish and later destroy a savepoint:

BEGIN;
    INSERT INTO table1 VALUES (3);
    SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (4);
    RELEASE SAVEPOINT my_savepoint;
COMMIT;

The above transaction will insert both 3 and 4.

Compatibility

This command conforms to the SQL standard. The standard specifies that the key word SAVEPOINT is mandatory, but PostgreSQL allows it to be omitted.
See Also

BEGIN, COMMIT, ROLLBACK, ROLLBACK TO SAVEPOINT, SAVEPOINT
RESET

RESET — restore the value of a run-time parameter to the default value

Synopsis

RESET configuration_parameter
RESET ALL

Description

RESET restores run-time parameters to their default values. RESET is an alternative spelling for

SET configuration_parameter TO DEFAULT

Refer to SET for details.

The default value is defined as the value that the parameter would have had, if no SET had ever been
issued for it in the current session. The actual source of this value might be a compiled-in default, the
configuration file, command-line options, or per-database or per-user default settings. This is subtly
different from defining it as “the value that the parameter had at session start”, because if the value
came from the configuration file, it will be reset to whatever is specified by the configuration file now.
See Chapter 19 for details.

The transactional behavior of RESET is the same as SET: its effects will be undone by transaction
rollback.

Parameters

configuration_parameter

Name of a settable run-time parameter. Available parameters are documented in Chapter 19 and
on the SET reference page.

ALL

Resets all settable run-time parameters to default values.

Examples

Set the timezone configuration variable to its default value:

RESET timezone;

Compatibility

RESET is a PostgreSQL extension.

See Also

SET, SHOW
REVOKE

REVOKE — remove access privileges

Synopsis

REVOKE [ GRANT OPTION FOR ]
  ( { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES | TRIGGER }
    [, ...] | ALL [ PRIVILEGES ] )
ON { [ TABLE ] table_name [, ...]
    | ALL TABLES IN SCHEMA schema_name [, ...] }
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( { SELECT | INSERT | UPDATE | REFERENCES } ( column_name
    [, ...] )
    [, ...] | ALL [ PRIVILEGES ] ( column_name [, ...] ) )
ON [ TABLE ] table_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( { USAGE | SELECT | UPDATE } )
    [, ...] | ALL [ PRIVILEGES ] )
ON { SEQUENCE sequence_name [, ...]
    | ALL SEQUENCES IN SCHEMA schema_name [, ...] }
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( { CREATE | CONNECT | TEMPORARY | TEMP } [, ...] | ALL
    PRIVILEGES )
ON DATABASE database_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( USAGE | ALL [ PRIVILEGES ] )
ON DOMAIN domain_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( USAGE | ALL [ PRIVILEGES ] )
ON FOREIGN DATA WRAPPER fdw_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  ( USAGE | ALL [ PRIVILEGES ] )
ON FOREIGN SERVER server_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
REVOKE [ GRANT OPTION FOR ]
  { EXECUTE | ALL [ PRIVILEGES ] }
  ON { { FUNCTION | PROCEDURE | ROUTINE } function_name
    [ ( [ [ argmode ] [ arg_name ] arg_type [, ...] ] ) [, ...] ]
    | ALL { FUNCTIONS | PROCEDURES | ROUTINES } }
  SCHEMA schema_name [, ...] }
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { USAGE | ALL [ PRIVILEGES ] }
  ON LANGUAGE lang_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { { SELECT | UPDATE } [, ...] | ALL [ PRIVILEGES ] }
  ON LARGE OBJECT loid [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { { CREATE | USAGE } [, ...] | ALL [ PRIVILEGES ] }
  ON SCHEMA schema_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { CREATE | USAGE } [, ...]
  ON TABLESPACE tablespace_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { USAGE | ALL [ PRIVILEGES ] }
  ON TYPE type_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ ADMIN OPTION FOR ]
  role_name [, ...] FROM role_name [, ...]
  [ CASCADE | RESTRICT ]

Description

The REVOKE command revokes previously granted privileges from one or more roles. The key word PUBLIC refers to the implicitly defined group of all roles.

See the description of the GRANT command for the meaning of the privilege types.

Note that any particular role will have the sum of privileges granted directly to it, privileges granted to any role it is presently a member of, and privileges granted to PUBLIC. Thus, for example, revoking SELECT privilege from PUBLIC does not necessarily mean that all roles have lost SELECT privilege on the object: those who have it granted directly or via another role will still have it. Similarly, revoking SELECT from a user might not prevent that user from using SELECT if PUBLIC or another membership role still has SELECT rights.
If GRANT OPTION FOR is specified, only the grant option for the privilege is revoked, not the privilege itself. Otherwise, both the privilege and the grant option are revoked.

If a user holds a privilege with grant option and has granted it to other users then the privileges held by those other users are called dependent privileges. If the privilege or the grant option held by the first user is being revoked and dependent privileges exist, those dependent privileges are also revoked if CASCADE is specified; if it is not, the revoke action will fail. This recursive revocation only affects privileges that were granted through a chain of users that is traceable to the user that is the subject of this REVOKE command. Thus, the affected users might effectively keep the privilege if it was also granted through other users.

When revoking privileges on a table, the corresponding column privileges (if any) are automatically revoked on each column of the table, as well. On the other hand, if a role has been granted privileges on a table, then revoking the privileges from a table will have no effect.

When revoking membership in a role, GRANT OPTION is instead called ADMIN OPTION, but the behavior is similar. Note also that this form of the command does not allow the noise word GROUP.

Notes

Use psql's \dp command to display the privileges granted on existing tables and columns. See GRANT for information about the format. For non-table objects there are other \d commands that can display their privileges.

A user can only revoke privileges that were granted directly by that user. If, for example, user A has granted a privilege with grant option to user B, and user B has in turn granted it to user C, then user A cannot revoke the privilege directly from C. Instead, user A could revoke the grant option from user B and use the CASCADE option so that the privilege is in turn revoked from user C. For another example, if both A and B have granted the same privilege to C, A can revoke their own grant but not B’s grant, so C will still effectively have the privilege.

When a non-owner of an object attempts to REVOKE privileges on the object, the command will fail outright if the user has no privileges whatsoever on the object. As long as some privilege is available, the command will proceed, but it will revoke only those privileges for which the user has grant options. The REVOKE ALL PRIVILEGES forms will issue a warning message if no grant options are held, while the other forms will issue a warning if grant options for any of the privileges specifically named in the command are not held. (In principle these statements apply to the object owner as well, but since the owner is always treated as holding all grant options, the cases can never occur.)

If a superuser chooses to issue a GRANT or REVOKE command, the command is performed as though it were issued by the owner of the affected object. Since all privileges ultimately come from the object owner (possibly indirectly via chains of grant options), it is possible for a superuser to revoke all privileges, but this might require use of CASCADE as stated above.

REVOKE can also be done by a role that is not the owner of the affected object, but is a member of the role that owns the object, or is a member of a role that holds privileges WITH GRANT OPTION on the object. In this case the command is performed as though it were issued by the containing role that actually owns the object or holds the privileges WITH GRANT OPTION. For example, if table t1 is owned by role g1, of which role u1 is a member, then u1 can revoke privileges on t1 that are recorded as being granted by g1. This would include grants made by u1 as well as by other members of role g1.

If the role executing REVOKE holds privileges indirectly via more than one role membership path, it is unspecified which containing role will be used to perform the command. In such cases it is best practice to use SET ROLE to become the specific role you want to do the REVOKE as. Failure to do so might lead to revoking privileges other than the ones you intended, or not revoking anything at all.

Examples

Revoke insert privilege for the public on table films:
REVOKE INSERT ON films FROM PUBLIC;

Revoke all privileges from user manuel on view kinds:

REVOKE ALL PRIVILEGES ON kinds FROM manuel;

Note that this actually means "revoke all privileges that I granted".

Revoke membership in role admins from user joe:

REVOKE admins FROM joe;

**Compatibility**

The compatibility notes of the GRANT command apply analogously to REVOKE. The keyword RESTRICT or CASCADE is required according to the standard, but PostgreSQL assumes RESTRICT by default.

**See Also**

GRANT
ROLLBACK

ROLLBACK — abort the current transaction

Synopsis

ROLLBACK [ WORK | TRANSACTION ]

Description

ROLLBACK rolls back the current transaction and causes all the updates made by the transaction to be discarded.

Parameters

WORK
TRANSACTION

Optional key words. They have no effect.

Notes

Use COMMIT to successfully terminate a transaction.

Issuing ROLLBACK outside of a transaction block emits a warning and otherwise has no effect.

Examples

To abort all changes:

ROLLBACK;

Compatibility

The SQL standard only specifies the two forms ROLLBACK and ROLLBACK WORK. Otherwise, this command is fully conforming.

See Also

BEGIN, COMMIT, ROLLBACK TO SAVEPOINT
ROLLBACK PREPARED

ROLLBACK PREPARED — cancel a transaction that was earlier prepared for two-phase commit

Synopsis

```
ROLLBACK PREPARED transaction_id
```

Description

ROLLBACK PREPARED rolls back a transaction that is in prepared state.

Parameters

```
transaction_id
```

The transaction identifier of the transaction that is to be rolled back.

Notes

To roll back a prepared transaction, you must be either the same user that executed the transaction originally, or a superuser. But you do not have to be in the same session that executed the transaction.

This command cannot be executed inside a transaction block. The prepared transaction is rolled back immediately.

All currently available prepared transactions are listed in the `pg_prepared_xacts` system view.

Examples

Roll back the transaction identified by the transaction identifier `foobar`:

```
ROLLBACK PREPARED 'foobar';
```

Compatibility

ROLLBACK PREPARED is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.

See Also

PREPARE TRANSACTION, COMMIT PREPARED
ROLLBACK TO SAVEPOINT

ROLLBACK TO SAVEPOINT — roll back to a savepoint

Synopsis

ROLLBACK [ WORK | TRANSACTION ] TO [ SAVEPOINT ] savepoint_name

Description

Roll back all commands that were executed after the savepoint was established. The savepoint remains valid and can be rolled back to again later, if needed.

ROLLBACK TO SAVEPOINT implicitly destroys all savepoints that were established after the named savepoint.

Parameters

savepoint_name

The savepoint to roll back to.

Notes

Use RELEASE SAVEPOINT to destroy a savepoint without discarding the effects of commands executed after it was established.

Specifying a savepoint name that has not been established is an error.

Cursors have somewhat non-transactional behavior with respect to savepoints. Any cursor that is opened inside a savepoint will be closed when the savepoint is rolled back. If a previously opened cursor is affected by a FETCH or MOVE command inside a savepoint that is later rolled back, the cursor remains at the position that FETCH left it pointing to (that is, the cursor motion caused by FETCH is not rolled back). Closing a cursor is not undone by rolling back, either. However, other side-effects caused by the cursor’s query (such as side-effects of volatile functions called by the query) are rolled back if they occur during a savepoint that is later rolled back. A cursor whose execution causes a transaction to abort is put in a cannot-execute state, so while the transaction can be restored using ROLLBACK TO SAVEPOINT, the cursor can no longer be used.

Examples

To undo the effects of the commands executed after my_savepoint was established:

ROLLBACK TO SAVEPOINT my_savepoint;

Cursor positions are not affected by savepoint rollback:

BEGIN;

DECLARE foo CURSOR FOR SELECT 1 UNION SELECT 2;

SAVEPOINT foo;
ROLLBACK TO SAVEPOINT

FETCH 1 FROM foo;
?column?
--------
 1

ROLLBACK TO SAVEPOINT foo;

FETCH 1 FROM foo;
?column?
--------
 2

COMMIT;

Compatibility

The SQL standard specifies that the key word SAVEPOINT is mandatory, but PostgreSQL and Oracle allow it to be omitted. SQL allows only WORK, not TRANSACTION, as a noise word after ROLLBACK. Also, SQL has an optional clause AND [ NO ] CHAIN which is not currently supported by PostgreSQL. Otherwise, this command conforms to the SQL standard.

See Also

BEGIN, COMMIT, RELEASE SAVEPOINT, ROLLBACK, SAVEPOINT
SAVEPOINT

SAVEPOINT — define a new savepoint within the current transaction

Synopsis

SAVEPOINT savepoint_name

Description

SAVEPOINT establishes a new savepoint within the current transaction.

A savepoint is a special mark inside a transaction that allows all commands that are executed after it was established to be rolled back, restoring the transaction state to what it was at the time of the savepoint.

Parameters

savepoint_name

The name to give to the new savepoint.

Notes

Use ROLLBACK TO SAVEPOINT to rollback to a savepoint. Use RELEASE SAVEPOINT to destroy a savepoint, keeping the effects of commands executed after it was established.

Savepoints can only be established when inside a transaction block. There can be multiple savepoints defined within a transaction.

Examples

To establish a savepoint and later undo the effects of all commands executed after it was established:

BEGIN;
    INSERT INTO table1 VALUES (1);
    SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (2);
    ROLLBACK TO SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (3);
COMMIT;

The above transaction will insert the values 1 and 3, but not 2.

To establish and later destroy a savepoint:

BEGIN;
    INSERT INTO table1 VALUES (3);
    SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (4);
    RELEASE SAVEPOINT my_savepoint;
COMMIT;
The above transaction will insert both 3 and 4.

**Compatibility**

SQL requires a savepoint to be destroyed automatically when another savepoint with the same name is established. In PostgreSQL, the old savepoint is kept, though only the more recent one will be used when rolling back or releasing. (Releasing the newer savepoint with `RELEASE SAVEPOINT` will cause the older one to again become accessible to `ROLLBACK TO SAVEPOINT` and `RELEASE SAVEPOINT`.) Otherwise, `SAVEPOINT` is fully SQL conforming.

**See Also**

BEGIN, COMMIT, RELEASE SAVEPOINT, ROLLBACK, ROLLBACK TO SAVEPOINT
SECURITY LABEL

SECURITY LABEL — define or change a security label applied to an object

Synopsis

SECURITY LABEL [ FOR provider ] ON
{
  TABLE object_name |
  COLUMN table_name.column_name |
  AGGREGATE aggregate_name ( aggregate_signature ) |
  DATABASE object_name |
  DOMAIN object_name |
  EVENT TRIGGER object_name |
  FOREIGN TABLE object_name
  FUNCTION function_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  LARGE OBJECT large_object_oid |
  MATERIALIZED VIEW object_name |
  [ PROCEDURAL ] LANGUAGE object_name |
  PROCEDURE procedure_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  PUBLICATION object_name |
  ROLE object_name |
  ROUTINE routine_name [ ( [ [ argmode ] [ argname ] argtype [, ...] ] ) ] |
  SCHEMA object_name |
  SEQUENCE object_name |
  SUBSCRIPTION object_name |
  TABLESPACE object_name |
  TYPE object_name |
  VIEW object_name
} IS 'label'

where aggregate_signature is:

* |
[ argmode ] [ argname ] argtype [, ... ] |
[ [ argmode ] [ argname ] argtype [, ... ] ] ORDER BY [ argmode ] |
[ argname ] argtype [, ... ]

Description

SECURITY LABEL applies a security label to a database object. An arbitrary number of security labels, one per label provider, can be associated with a given database object. Label providers are loadable modules which register themselves by using the function register_label_provider.

Note

register_label_provider is not an SQL function; it can only be called from C code loaded into the backend.
The label provider determines whether a given label is valid and whether it is permissible to assign that label to a given object. The meaning of a given label is likewise at the discretion of the label provider. PostgreSQL places no restrictions on whether or how a label provider must interpret security labels; it merely provides a mechanism for storing them. In practice, this facility is intended to allow integration with label-based mandatory access control (MAC) systems such as SE-Linux. Such systems make all access control decisions based on object labels, rather than traditional discretionary access control (DAC) concepts such as users and groups.

Parameters

object_name
table_name.column_name
aggregate_name
function_name
procedure_name
routine_name

The name of the object to be labeled. Names of tables, aggregates, domains, foreign tables, functions, procedures, routines, sequences, types, and views can be schema-qualified.

provider

The name of the provider with which this label is to be associated. The named provider must be loaded and must consent to the proposed labeling operation. If exactly one provider is loaded, the provider name may be omitted for brevity.

argmode

The mode of a function, procedure, or aggregate argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that SECURITY LABEL does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function's identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname

The name of a function, procedure, or aggregate argument. Note that SECURITY LABEL does not actually pay any attention to argument names, since only the argument data types are needed to determine the function's identity.

argtype

The data type of a function, procedure, or aggregate argument.

large_object_oid

The OID of the large object.

PROCEDURAL

This is a noise word.

label

The new security label, written as a string literal; or NULL to drop the security label.

Examples

The following example shows how the security label of a table might be changed.
SECURITY LABEL FOR selinux ON TABLE mytable IS 'system_u:object_r:sepgsql_table_t:s0';

Compatibility

There is no SECURITY LABEL command in the SQL standard.

See Also

sepgsql/src/test/modules/dummy_seclabel
SELECT

SELECT, TABLE, WITH — retrieve rows from a table or view

Synopsis

[ WITH [ RECURSIVE ] with_query [, ... ] ]
SELECT [ ALL | DISTINCT [ ON ( expression [, ... ] ) ] ]
[ * | expression [ [ AS ] output_name ] [, ... ] ]
FROM from_item [, ... ]
WHERE condition
GROUP BY grouping_element [, ... ]
HAVING condition [, ... ]
WINDOW window_name AS ( window_definition ) [, ... ]
( UNION | INTERSECT | EXCEPT ) [ ALL | DISTINCT ] select
ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS
( FIRST | LAST ) ] [, ... ]
LIMIT { count | ALL }
OFFSET start [ ROW | ROWS ]
FETCH { FIRST | NEXT } { count | ROW | ROWS } ONLY
FOR { UPDATE | NO KEY UPDATE | SHARE | KEY SHARE }
OF table_name [, ... ] [ NOWAIT | SKIP LOCKED ] [...] ]

where from_item can be one of:

[ ONLY ] table_name [ * ] [ [ AS ] alias [ ( column_alias
[, ... ] ) ] ]
[ TABLESAMPLE sampling_method ( argument [, ... ] ) ]
[ REPEATABLE ( seed ) ]
[ LATERAL ] ( select ) [ AS ] alias [ ( column_alias
[, ... ] ) ]
with_query_name [ [ AS ] alias [ ( column_alias [, ... ] ) ] ]
[ LATERAL ] function_name ( [ argument [, ... ] ] )
[ WITH ORDINALITY ] [ [ AS ] alias [ ( column_alias
[, ... ] ) ] ]
[ LATERAL ] function_name ( [ argument [, ... ] ] ) [ AS ] alias
( column_definition [, ... ] )
[ LATERAL ] function_name ( [ argument [, ... ] ] ) AS
( column_definition [, ... ] )
[ LATERAL ] ROWS FROM( function_name ( [ argument [, ... ] ]
) [ AS ( column_definition [, ... ] ) ] [, ... ]
[ WITH ORDINALITY ] [ [ AS ] alias [ ( column_alias
[, ... ] ) ] ]
from_item [ NATURAL ] join_type from_item [ ON join_condition
USING ( join_column [, ... ] ) ]

and grouping_element can be one of:

( )
expression
( expression [, ... ] )
ROLLUP ( [ expression | ( expression [, ... ] ) ] [, ... ] )
CUBE ( [ expression | ( expression [, ... ] ) ] [, ... ] )
GROUPING SETS ( grouping_element [, ... ] )

1764
and with_query is:

```sql
WITH
    with_query_name [ ( column_name [, ...] ) ] AS ( select
            | values | insert | update | delete )
```

TABLE [ ONLY ] table_name [ * ]

**Description**

SELECT retrieves rows from zero or more tables. The general processing of SELECT is as follows:

1. All queries in the WITH list are computed. These effectively serve as temporary tables that can be referenced in the FROM list. A WITH query that is referenced more than once in FROM is computed only once. (See WITH Clause below.)

2. All elements in the FROM list are computed. (Each element in the FROM list is a real or virtual table.) If more than one element is specified in the FROM list, they are cross-joined together. (See FROM Clause below.)

3. If the WHERE clause is specified, all rows that do not satisfy the condition are eliminated from the output. (See WHERE Clause below.)

4. If the GROUP BY clause is specified, or if there are aggregate function calls, the output is combined into groups of rows that match on one or more values, and the results of aggregate functions are computed. If the HAVING clause is present, it eliminates groups that do not satisfy the given condition. (See GROUP BY Clause and HAVING Clause below.)

5. The actual output rows are computed using the SELECT output expressions for each selected row or row group. (See SELECT List below.)

6. SELECT DISTINCT eliminates duplicate rows from the result. SELECT DISTINCT ON eliminates rows that match on all the specified expressions. SELECT ALL (the default) will return all candidate rows, including duplicates. (See DISTINCT Clause below.)

7. Using the operators UNION, INTERSECT, and EXCEPT, the output of more than one SELECT statement can be combined to form a single result set. The UNION operator returns all rows that are in one or both of the result sets. The INTERSECT operator returns all rows that are strictly in both result sets. The EXCEPT operator returns the rows that are in the first result set but not in the second. In all three cases, duplicate rows are eliminated unless ALL is specified. The noise word DISTINCT can be added to explicitly specify eliminating duplicate rows. Notice that DISTINCT is the default behavior here, even though ALL is the default for SELECT itself. (See UNION Clause, INTERSECT Clause, and EXCEPT Clause below.)

8. If the ORDER BY clause is specified, the returned rows are sorted in the specified order. If ORDER BY is not given, the rows are returned in whatever order the system finds fastest to produce. (See ORDER BY Clause below.)

9. If the LIMIT (or FETCH FIRST) or OFFSET clause is specified, the SELECT statement only returns a subset of the result rows. (See LIMIT Clause below.)

10. If FOR UPDATE, FOR NO KEY UPDATE, FOR SHARE or FOR KEY SHARE is specified, the SELECT statement locks the selected rows against concurrent updates. (See The Locking Clause below.)

You must have SELECT privilege on each column used in a SELECT command. The use of FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE or FOR KEY SHARE requires UPDATE privilege as well (for at least one column of each table so selected).
Parameters

**WITH Clause**

The **WITH** clause allows you to specify one or more subqueries that can be referenced by name in the primary query. The subqueries effectively act as temporary tables or views for the duration of the primary query. Each subquery can be a **SELECT**, **TABLE**, **VALUES**, **INSERT**, **UPDATE** or **DELETE** statement. When writing a data-modifying statement (**INSERT, UPDATE or DELETE**) in **WITH**, it is usual to include a **RETURNING** clause. It is the output of **RETURNING**, not the underlying table that the statement modifies, that forms the temporary table that is read by the primary query. If **RETURNING** is omitted, the statement is still executed, but it produces no output so it cannot be referenced as a table by the primary query.

A name (without schema qualification) must be specified for each **WITH** query. Optionally, a list of column names can be specified; if this is omitted, the column names are inferred from the subquery.

If **RECURSIVE** is specified, it allows a **SELECT** subquery to reference itself by name. Such a subquery must have the form

```
non_recursive_term UNION [ ALL | DISTINCT ] recursive_term
```

where the recursive self-reference must appear on the right-hand side of the **UNION**. Only one recursive self-reference is permitted per query. Recursive data-modifying statements are not supported, but you can use the results of a recursive **SELECT** query in a data-modifying statement. See Section 7.8 for an example.

Another effect of **RECURSIVE** is that **WITH** queries need not be ordered: a query can reference another one that is later in the list. (However, circular references, or mutual recursion, are not implemented.) Without **RECURSIVE**, **WITH** queries can only reference sibling **WITH** queries that are earlier in the **WITH** list.

A key property of **WITH** queries is that they are evaluated only once per execution of the primary query, even if the primary query refers to them more than once. In particular, data-modifying statements are guaranteed to be executed once and only once, regardless of whether the primary query reads all or any of their output.

The primary query and the **WITH** queries are all (notionally) executed at the same time. This implies that the effects of a data-modifying statement in **WITH** cannot be seen from other parts of the query, other than by reading its **RETURNING** output. If two such data-modifying statements attempt to modify the same row, the results are unspecified.

See Section 7.8 for additional information.

**FROM Clause**

The **FROM** clause specifies one or more source tables for the **SELECT**. If multiple sources are specified, the result is the Cartesian product (cross join) of all the sources. But usually qualification conditions are added (via **WHERE**) to restrict the returned rows to a small subset of the Cartesian product.

The **FROM** clause can contain the following elements:

```
table_name
```

The name (optionally schema-qualified) of an existing table or view. If **ONLY** is specified before the table name, only that table is scanned. If **ONLY** is not specified, the table and all its descendant tables (if any) are scanned. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.
**alias**

A substitute name for the FROM item containing the alias. An alias is used for brevity or to eliminate ambiguity for self-joins (where the same table is scanned multiple times). When an alias is provided, it completely hides the actual name of the table or function; for example given `FROM foo AS f`, the remainder of the SELECT must refer to this FROM item as `f` not `foo`. If an alias is written, a column alias list can also be written to provide substitute names for one or more columns of the table.

**TABLESAMPLE** `sampling_method` ( `argument` [, ...] ) [ REPEATABLE ( `seed` ) ]

A TABLESAMPLE clause after a `table_name` indicates that the specified `sampling_method` should be used to retrieve a subset of the rows in that table. This sampling precedes the application of any other filters such as WHERE clauses. The standard PostgreSQL distribution includes two sampling methods, `BERNOULLI` and `SYSTEM`, and other sampling methods can be installed in the database via extensions.

The `BERNOULLI` and `SYSTEM` sampling methods each accept a single `argument` which is the fraction of the table to sample, expressed as a percentage between 0 and 100. This argument can be any real-valued expression. (Other sampling methods might accept more or different arguments.) These two methods each return a randomly-chosen sample of the table that will contain approximately the specified percentage of the table's rows. The `BERNOULLI` method scans the whole table and selects or ignores individual rows independently with the specified probability. The `SYSTEM` method does block-level sampling with each block having the specified chance of being selected; all rows in each selected block are returned. The `SYSTEM` method is significantly faster than the `BERNOULLI` method when small sampling percentages are specified, but it may return a less-random sample of the table as a result of clustering effects.

The optional `REPEATABLE` clause specifies a `seed` number or expression to use for generating random numbers within the sampling method. The seed value can be any non-null floating-point value. Two queries that specify the same seed and `argument` values will select the same sample of the table, if the table has not been changed meanwhile. But different seed values will usually produce different samples. If `REPEATABLE` is not given then a new random sample is selected for each query, based upon a system-generated seed. Note that some add-on sampling methods do not accept `REPEATABLE`, and will always produce new samples on each use.

**select**

A sub-SELECT can appear in the FROM clause. This acts as though its output were created as a temporary table for the duration of this single SELECT command. Note that the sub-SELECT must be surrounded by parentheses, and an alias must be provided for it. A VALUES command can also be used here.

**with_query_name**

A WITH query is referenced by writing its name, just as though the query's name were a table name. (In fact, the WITH query hides any real table of the same name for the purposes of the primary query. If necessary, you can refer to a real table of the same name by schema-qualifying the table's name.) An alias can be provided in the same way as for a table.

**function_name**

Function calls can appear in the FROM clause. (This is especially useful for functions that return result sets, but any function can be used.) This acts as though the function's output were created as a temporary table for the duration of this single SELECT command. When the optional `WITH ORDINALITY` clause is added to the function call, a new column is appended after all the function's output columns with numbering for each row.

An alias can be provided in the same way as for a table. If an alias is written, a column alias list can also be written to provide substitute names for one or more attributes of the function's composite return type, including the column added by `ORDINALITY` if present.
Multiple function calls can be combined into a single FROM-clause item by surrounding them with `ROWS FROM(...)` The output of such an item is the concatenation of the first row from each function, then the second row from each function, etc. If some of the functions produce fewer rows than others, null values are substituted for the missing data, so that the total number of rows returned is always the same as for the function that produced the most rows.

If the function has been defined as returning the `record` data type, then an alias or the keyword `AS` must be present, followed by a column definition list in the form `{column_name data_type [, ...]}`. The column definition list must match the actual number and types of columns returned by the function.

When using the `ROWS FROM(...)` syntax, if one of the functions requires a column definition list, it's preferred to put the column definition list after the function call inside `ROWS FROM(...)`. A column definition list can be placed after the `ROWS FROM(...)` construct only if there's just a single function and no `WITH ORDINALITY` clause.

To use `ORDINALITY` together with a column definition list, you must use the `ROWS FROM(...)` syntax and put the column definition list inside `ROWS FROM(...)`.

### join_type

One of

- `[ INNER ] JOIN`
- `LEFT [ OUTER ] JOIN`
- `RIGHT [ OUTER ] JOIN`
- `FULL [ OUTER ] JOIN`
- `CROSS JOIN`

For the `INNER` and `OUTER` join types, a join condition must be specified, namely exactly one of `NATURAL`, `ON join_condition`, or `USING (join_column [, ...])`. See below for the meaning. For `CROSS JOIN`, none of these clauses can appear.

A JOI N clause combines two FROM items, which for convenience we will refer to as “tables”, though in reality they can be any type of FROM item. Use parentheses if necessary to determine the order of nesting. In the absence of parentheses, JOINs nest left-to-right. In any case JOIN binds more tightly than the commas separating FROM-list items.

`CROSS JOIN` and `INNER JOIN` produce a simple Cartesian product, the same result as you get from listing the two tables at the top level of FROM, but restricted by the join condition (if any). `CROSS JOIN` is equivalent to `INNER JOIN ON (TRUE)`, that is, no rows are removed by qualification. These join types are just a notational convenience, since they do nothing you couldn't do with plain FROM and WHERE.

`LEFT OUTER JOIN` returns all rows in the qualified Cartesian product (i.e., all combined rows that pass its join condition), plus one copy of each row in the left-hand table for which there was no right-hand row that passed the join condition. This left-hand row is extended to the full width of the joined table by inserting null values for the right-hand columns. Note that only the JOIN clause's own condition is considered while deciding which rows have matches. Outer conditions are applied afterwards.

Conversely, `RIGHT OUTER JOIN` returns all the joined rows, plus one row for each unmatched right-hand row (extended with nulls on the left). This is just a notational convenience, since you could convert it to a `LEFT OUTER JOIN` by switching the left and right tables.

`FULL OUTER JOIN` returns all the joined rows, plus one row for each unmatched left-hand row (extended with nulls on the right), plus one row for each unmatched right-hand row (extended with nulls on the left).
ON join_condition

join_condition is an expression resulting in a value of type boolean (similar to a WHERE clause) that specifies which rows in a join are considered to match.

USING ( join_column [, ...] )

A clause of the form USING ( a, b, ... ) is shorthand for ON left_table.a = right_table.a AND left_table.b = right_table.b .... Also, USING implies that only one of each pair of equivalent columns will be included in the join output, not both.

NATURAL

NATURAL is shorthand for a USING list that mentions all columns in the two tables that have matching names. If there are no common column names, NATURAL is equivalent to ON TRUE.

LATERAL

The LATERAL key word can precede a sub-SELECT FROM item. This allows the sub-SELECT to refer to columns of FROM items that appear before it in the FROM list. (Without LATERAL, each sub-SELECT is evaluated independently and so cannot cross-reference any other FROM item.)

LATERAL can also precede a function-call FROM item, but in this case it is a noise word, because the function expression can refer to earlier FROM items in any case.

A LATERAL item can appear at top level in the FROM list, or within a JOIN tree. In the latter case it can also refer to any items that are on the left-hand side of a JOIN that it is on the right-hand side of.

When a FROM item contains LATERAL cross-references, evaluation proceeds as follows: for each row of the FROM item providing the cross-referenced column(s), or set of rows of multiple FROM items providing the columns, the LATERAL item is evaluated using that row or row set's values of the columns. The resulting row(s) are joined as usual with the rows they were computed from. This is repeated for each row or set of rows from the column source table(s).

The column source table(s) must be INNER or LEFT joined to the LATERAL item, else there would not be a well-defined set of rows from which to compute each set of rows for the LATERAL item. Thus, although a construct such as X RIGHT JOIN LATERAL Y is syntactically valid, it is not actually allowed for Y to reference X.

WHERE Clause

The optional WHERE clause has the general form

WHERE condition

where condition is any expression that evaluates to a result of type boolean. Any row that does not satisfy this condition will be eliminated from the output. A row satisfies the condition if it returns true when the actual row values are substituted for any variable references.

GROUP BY Clause

The optional GROUP BY clause has the general form

GROUP BY grouping_element [, ...]

GROUP BY will condense into a single row all selected rows that share the same values for the grouped expressions. An expression used inside a grouping_element can be an input column name, or the name or ordinal number of an output column (SELECT list item), or an arbitrary expression
formed from input-column values. In case of ambiguity, a GROUP BY name will be interpreted as an input-column name rather than an output column name.

If any of GROUPING SETS, ROLLUP or CUBE are present as grouping elements, then the GROUP BY clause as a whole defines some number of independent grouping sets. The effect of this is equivalent to constructing a UNION ALL between subqueries with the individual grouping sets as their GROUP BY clauses. For further details on the handling of grouping sets see Section 7.2.4.

Aggregate functions, if any are used, are computed across all rows making up each group, producing a separate value for each group. (If there are aggregate functions but no GROUP BY clause, the query is treated as having a single group comprising all the selected rows.) The set of rows fed to each aggregate function can be further filtered by attaching a FILTER clause to the aggregate function call; see Section 4.2.7 for more information. When a FILTER clause is present, only those rows matching it are included in the input to that aggregate function.

When GROUP BY is present, or any aggregate functions are present, it is not valid for the SELECT list expressions to refer to ungrouped columns except within aggregate functions or when the ungrouped column is functionally dependent on the grouped columns, since there would otherwise be more than one possible value to return for an ungrouped column. A functional dependency exists if the grouped columns (or a subset thereof) are the primary key of the table containing the ungrouped column.

Keep in mind that all aggregate functions are evaluated before evaluating any “scalar” expressions in the HAVING clause or SELECT list. This means that, for example, a CASE expression cannot be used to skip evaluation of an aggregate function; see Section 4.2.14.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with GROUP BY.

**HAVING Clause**

The optional HAVING clause has the general form

```
HAVING condition
```

where condition is the same as specified for the WHERE clause.

HAVING eliminates group rows that do not satisfy the condition. HAVING is different from WHERE: WHERE filters individual rows before the application of GROUP BY, while HAVING filters group rows created by GROUP BY. Each column referenced in condition must unambiguously reference a grouping column, unless the reference appears within an aggregate function or the ungrouped column is functionally dependent on the grouping columns.

The presence of HAVING turns a query into a grouped query even if there is no GROUP BY clause. This is the same as what happens when the query contains aggregate functions but no GROUP BY clause. All the selected rows are considered to form a single group, and the SELECT list and HAVING clause can only reference table columns from within aggregate functions. Such a query will emit a single row if the HAVING condition is true, zero rows if it is not true.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with HAVING.

**WINDOW Clause**

The optional WINDOW clause has the general form

```
WINDOW window_name AS ( window_definition ) [, ...]
```

where window_name is a name that can be referenced from OVER clauses or subsequent window definitions, and window_definition is
SELECT

[ existing_window_name ]
[ PARTITION BY expression [, ...] ]
[ ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS
  ( FIRST | LAST ) ] [, ...] ]
[ frame_clause ]

If an existing_window_name is specified it must refer to an earlier entry in the WINDOW list; the new window copies its partitioning clause from that entry, as well as its ordering clause if any. In this case the new window cannot specify its own PARTITION BY clause, and it can specify ORDER BY only if the copied window does not have one. The new window always uses its own frame clause; the copied window must not specify a frame clause.

The elements of the PARTITION BY list are interpreted in much the same fashion as elements of a GROUP BY Clause, except that they are always simple expressions and never the name or number of an output column. Another difference is that these expressions can contain aggregate function calls, which are not allowed in a regular GROUP BY clause. They are allowed here because windowing occurs after grouping and aggregation.

Similarly, the elements of the ORDER BY list are interpreted in much the same fashion as elements of an ORDER BY Clause, except that the expressions are always taken as simple expressions and never the name or number of an output column.

The optional frame_clause defines the window frame for window functions that depend on the frame (not all do). The window frame is a set of related rows for each row of the query (called the current row). The frame_clause can be one of

{ RANGE | ROWS | GROUPS } frame_start [ frame_exclusion ]
{ RANGE | ROWS | GROUPS } BETWEEN frame_start AND frame_end
[ frame_exclusion ]

where frame_start and frame_end can be one of

UNBOUNDED PRECEDING
offset PRECEDING
CURRENT ROW
offset FOLLOWING
UNBOUNDED FOLLOWING

and frame_exclusion can be one of

EXCLUDE CURRENT ROW
EXCLUDE GROUP
EXCLUDE TIES
EXCLUDE NO OTHERS

If frame_end is omitted it defaults to CURRENT ROW. Restrictions are that frame_start cannot be UNBOUNDED FOLLOWING, frame_end cannot be UNBOUNDED PRECEDING, and the frame_end choice cannot appear earlier in the above list of frame_start and frame_end options than the frame_start choice does — for example RANGE BETWEEN CURRENT ROW AND offset PRECEDING is not allowed.

The default framing option is RANGE UNBOUNDED PRECEDING, which is the same as RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW; it sets the frame to be all rows from the partition start up through the current row’s last peer (a row that the window’s ORDER BY clause considers equivalent to the current row; all rows are peers if there is no ORDER BY). In general, UNBOUNDED PRECEDING means that the frame starts with the first row of the partition, and similarly
UNBOUNDED FOLLOWING means that the frame ends with the last row of the partition, regardless of RANGE, ROWS or GROUPS mode. In ROWS mode, CURRENT ROW means that the frame starts or ends with the current row; but in RANGE or GROUPS mode it means that the frame starts or ends with the current row’s first or last peer in the ORDER BY ordering. The offset PRECEDING and offset FOLLOWING options vary in meaning depending on the frame mode. In ROWS mode, the offset is an integer indicating that the frame starts or ends that many rows before or after the current row. In GROUPS mode, the offset is an integer indicating that the frame starts or ends that many peer groups before or after the current row’s peer group, where a peer group is a group of rows that are equivalent according to the window’s ORDER BY clause. In RANGE mode, use of an offset option requires that there be exactly one ORDER BY column in the window definition. Then the frame contains those rows whose ordering column value is no more than offset less than (for PRECEDING) or more than (for FOLLOWING) the current row’s ordering column value. In these cases the data type of the offset expression depends on the data type of the ordering column. For numeric ordering columns it is typically of the same type as the ordering column, but for datetime ordering columns it is an interval. In all these cases, the value of the offset must be non-null and non-negative. Also, while the offset does not have to be a simple constant, it cannot contain variables, aggregate functions, or window functions.

The frame_exclusion option allows rows around the current row to be excluded from the frame, even if they would be included according to the frame start and frame end options. EXCLUDE CURRENT ROW excludes the current row from the frame. EXCLUDE GROUP excludes the current row and its ordering peers from the frame. EXCLUDE TIES excludes any peers of the current row from the frame, but not the current row itself. EXCLUDE NO OTHERS simply specifies explicitly the default behavior of not excluding the current row or its peers.

Beware that the ROWS mode can produce unpredictable results if the ORDER BY ordering does not order the rows uniquely. The RANGE and GROUPS modes are designed to ensure that rows that are peers in the ORDER BY ordering are treated alike: all rows of a given peer group will be in the frame or excluded from it.

The purpose of a WINDOW clause is to specify the behavior of window functions appearing in the query’s SELECT List or ORDER BY Clause. These functions can reference the WINDOW clause entries by name in their OVER clauses. A WINDOW clause entry does not have to be referenced anywhere, however; if it is not used in the query it is simply ignored. It is possible to use window functions without any WINDOW clause at all, since a window function call can specify its window definition directly in its OVER clause. However, the WINDOW clause saves typing when the same window definition is needed for more than one window function.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with WINDOW.

Window functions are described in detail in Section 3.5, Section 4.2.8, and Section 7.2.5.

**SELECT List**

The SELECT list (between the key words SELECT and FROM) specifies expressions that form the output rows of the SELECT statement. The expressions can (and usually do) refer to columns computed in the FROM clause.

Just as in a table, every output column of a SELECT has a name. In a simple SELECT this name is just used to label the column for display, but when the SELECT is a sub-query of a larger query, the name is seen by the larger query as the column name of the virtual table produced by the sub-query. To specify the name to use for an output column, write AS output_name after the column’s expression. (You can omit AS, but only if the desired output name does not match any PostgreSQL keyword (see Appendix C). For protection against possible future keyword additions, it is recommended that you always either write AS or double-quote the output name.) If you do not specify a column name, a name is chosen automatically by PostgreSQL. If the column’s expression is a simple column reference then the chosen name is the same as that column’s name. In more complex cases a function or type name may be used, or the system may fall back on a generated name such as ?column?.
An output column's name can be used to refer to the column's value in ORDER BY and GROUP BY clauses, but not in the WHERE or HAVING clauses; there you must write out the expression instead.

Instead of an expression, * can be written in the output list as a shorthand for all the columns of the selected rows. Also, you can write table_name.* as a shorthand for the columns coming from just that table. In these cases it is not possible to specify new names with AS; the output column names will be the same as the table columns' names.

According to the SQL standard, the expressions in the output list should be computed before applying DISTINCT, ORDER BY, or LIMIT. This is obviously necessary when using DISTINCT, since otherwise it's not clear what values are being made distinct. However, in many cases it is convenient if output expressions are computed after ORDER BY and LIMIT; particularly if the output list contains any volatile or expensive functions. With that behavior, the order of function evaluations is more intuitive and there will not be evaluations corresponding to rows that never appear in the output. PostgreSQL will effectively evaluate output expressions after sorting and limiting, so long as those expressions are not referenced in DISTINCT, ORDER BY or GROUP BY. (As a counterexample, SELECT f(x) FROM tab ORDER BY 1 clearly must evaluate f(x) before sorting.) Output expressions that contain set-returning functions are effectively evaluated after sorting and before limiting, so that LIMIT will act to cut off the output from a set-returning function.

**Note**

PostgreSQL versions before 9.6 did not provide any guarantees about the timing of evaluation of output expressions versus sorting and limiting; it depended on the form of the chosen query plan.

**DISTINCT Clause**

If SELECT DISTINCT is specified, all duplicate rows are removed from the result set (one row is kept from each group of duplicates). SELECT ALL specifies the opposite: all rows are kept; that is the default.

SELECT DISTINCT ON (expression [, ...] ) keeps only the first row of each set of rows where the given expressions evaluate to equal. The DISTINCT ON expressions are interpreted using the same rules as for ORDER BY (see above). Note that the “first row” of each set is unpredictable unless ORDER BY is used to ensure that the desired row appears first. For example:

```
SELECT DISTINCT ON (location) location, time, report
FROM weather_reports
ORDER BY location, time DESC;
```

retrieves the most recent weather report for each location. But if we had not used ORDER BY to force descending order of time values for each location, we'd have gotten a report from an unpredictable time for each location.

The DISTINCT ON expression(s) must match the leftmost ORDER BY expression(s). The ORDER BY clause will normally contain additional expression(s) that determine the desired precedence of rows within each DISTINCT ON group.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with DISTINCT.

**UNION Clause**

The UNION clause has this general form:
**SELECT**

\[
\text{select_statement} \ \text{UNION} \ [ \ \text{ALL} \ | \ \text{DISTINCT} \ ] \ \text{select_statement}
\]

\text{select_statement} \ is \ any \ \text{SELECT} \ \text{statement} \ \text{without} \ \text{an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause.} \ (\text{ORDER BY and LIMIT can be} \ \text{attached to a subexpression} \ \text{if it is} \ \text{enclosed in parentheses. Without parentheses, these clauses will be taken} \ \text{to apply to the result of the UNION, not to its right-hand input expression.)}

The \text{UNION} \ \text{operator} \ \text{computes} \ \text{the set union} \ \text{of} \ \text{the rows returned} \ \text{by} \ \text{the involved} \ \text{SELECT statements.} \ \text{A row} \ \text{is} \ \text{in} \ \text{the set union} \ \text{of} \ \text{two result sets} \ \text{if} \ \text{it appears} \ \text{in at least one of the result sets.} \ \text{The two SELECT statements that represent} \ \text{the direct operands} \ \text{of} \ \text{the UNION must produce} \ \text{the same number of columns, and corresponding columns must be of compatible data types.}

The \text{result of UNION} \ \text{does not contain any duplicate rows unless} \ \text{the ALL option is specified.} \ \text{ALL prevents} \ \text{elimination} \ \text{of duplicates. (Therefore, UNION ALL is usually significantly quicker than UNION; use ALL when you can.) DISTINCT can be written to explicitly specify} \ \text{the default behavior of eliminating duplicate rows.}

\text{Multiple UNION operators} \ \text{in the same SELECT statement are evaluated} \ \text{left to right, unless otherwise indicated by parentheses.}

\text{Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified} \ \text{either for} \ \text{a UNION result} \ \text{or for any input of} \ \text{a UNION.}

**INTERSECT Clause**

The \text{INTERSECT clause} \ \text{has} \ \text{this general form:}

\[
\text{select_statement} \ \text{INTERSECT} \ [ \ \text{ALL} \ | \ \text{DISTINCT} \ ] \ \text{select_statement}
\]

\text{select_statement} \ is \ any \ \text{SELECT} \ \text{statement} \ \text{without} \ \text{an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause.}

The \text{INTERSECT} \ \text{operator} \ \text{computes} \ \text{the set intersection} \ \text{of} \ \text{the rows returned} \ \text{by} \ \text{the involved} \ \text{SELECT statements.} \ \text{A row} \ \text{is} \ \text{in} \ \text{the intersection} \ \text{of} \ \text{two result sets} \ \text{if} \ \text{it appears} \ \text{in both result sets.}

The \text{result} \ \text{of INTERSECT} \ \text{does not contain any duplicate rows unless the ALL option is specified.} \ \text{With ALL, a row} \ \text{that} \ \text{has} \ m \ \text{duplicates} \ \text{in the left table and} \ n \ \text{duplicates} \ \text{in the right table will appear} \ \min(m,n) \ \text{times} \ \text{in the result set. DISTINCT can be written to explicitly specify} \ \text{the default behavior of eliminating duplicate rows.}

\text{Multiple INTERSECT operators} \ \text{in the same SELECT statement are evaluated} \ \text{left to right, unless parentheses} \ \text{dictate otherwise. INTERSECT binds more tightly than UNION. That is, a UNION B INTERSECT C will be read as} \ \text{A UNION (B INTERSECT C).}

\text{Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified} \ \text{either for} \ \text{an INTERSECT result} \ \text{or for any input of an INTERSECT.}

**EXCEPT Clause**

The \text{EXCEPT clause} \ \text{has} \ \text{this general form:}

\[
\text{select_statement} \ \text{EXCEPT} \ [ \ \text{ALL} \ | \ \text{DISTINCT} \ ] \ \text{select_statement}
\]

\text{select_statement} \ is \ any \ \text{SELECT} \ \text{statement} \ \text{without} \ \text{an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause.}

The \text{EXCEPT} \ \text{operator} \ \text{computes} \ \text{the set of} \ \text{rows that are in the result} \ \text{of the left SELECT statement but not in the result of the right one.}
The result of EXCEPT does not contain any duplicate rows unless the ALL option is specified. With ALL, a row that has \( m \) duplicates in the left table and \( n \) duplicates in the right table will appear \( \max(m-n,0) \) times in the result set. DISTINCT can be written to explicitly specify the default behavior of eliminating duplicate rows.

Multiple EXCEPT operators in the same SELECT statement are evaluated left to right, unless parentheses dictate otherwise. EXCEPT binds at the same level as UNION.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified either for an EXCEPT result or for any input of an EXCEPT.

**ORDER BY Clause**

The optional ORDER BY clause has this general form:

```
ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ...]
```

The ORDER BY clause causes the result rows to be sorted according to the specified expression(s). If two rows are equal according to the leftmost expression, they are compared according to the next expression and so on. If they are equal according to all specified expressions, they are returned in an implementation-dependent order.

Each expression can be the name or ordinal number of an output column (SELECT list item), or it can be an arbitrary expression formed from input-column values.

The ordinal number refers to the ordinal (left-to-right) position of the output column. This feature makes it possible to define an ordering on the basis of a column that does not have a unique name. This is never absolutely necessary because it is always possible to assign a name to an output column using the AS clause.

It is also possible to use arbitrary expressions in the ORDER BY clause, including columns that do not appear in the SELECT output list. Thus the following statement is valid:

```
SELECT name FROM distributors ORDER BY code;
```

A limitation of this feature is that an ORDER BY clause applying to the result of a UNION, INTERSECT, or EXCEPT clause can only specify an output column name or number, not an expression.

If an ORDER BY expression is a simple name that matches both an output column name and an input column name, ORDER BY will interpret it as the output column name. This is the opposite of the choice that GROUP BY will make in the same situation. This inconsistency is made to be compatible with the SQL standard.

Optionally one can add the key word ASC (ascending) or DESC (descending) after any expression in the ORDER BY clause. If not specified, ASC is assumed by default. Alternatively, a specific ordering operator name can be specified in the USING clause. An ordering operator must be a less-than or greater-than member of some B-tree operator family. ASC is usually equivalent to USING < and DESC is usually equivalent to USING >. (But the creator of a user-defined data type can define exactly what the default sort ordering is, and it might correspond to operators with other names.)

If NULLS LAST is specified, null values sort after all non-null values; if NULLS FIRST is specified, null values sort before all non-null values. If neither is specified, the default behavior is NULLS LAST when ASC is specified or implied, and NULLS FIRST when DESC is specified (thus, the default is to act as though nulls are larger than non-nulls). When USING is specified, the default nulls ordering depends on whether the operator is a less-than or greater-than operator.

Note that ordering options apply only to the expression they follow; for example ORDER BY \( x, y \) DESC does not mean the same thing as ORDER BY \( x \) DESC, \( y \) DESC.
Character-string data is sorted according to the collation that applies to the column being sorted. That can be overridden at need by including a \texttt{COLLATE} clause in the \textit{expression}, for example \texttt{ORDER BY mycolumn COLLATE "en_US"}. For more information see Section 4.2.10 and Section 23.2.

\section*{LIMIT Clause}

The \texttt{LIMIT} clause consists of two independent sub-clauses:

\begin{verbatim}
LIMIT \{ count | ALL \}
OFFSET start
\end{verbatim}

\textit{count} specifies the maximum number of rows to return, while \textit{start} specifies the number of rows to skip before starting to return rows. When both are specified, \textit{start} rows are skipped before starting to count the \textit{count} rows to be returned.

If the \textit{count} expression evaluates to \texttt{NULL}, it is treated as \texttt{LIMIT ALL}, i.e., no limit. If \textit{start} evaluates to \texttt{NULL}, it is treated the same as \texttt{OFFSET 0}.

SQL:2008 introduced a different syntax to achieve the same result, which PostgreSQL also supports. It is:

\begin{verbatim}
OFFSET start \{ ROW | ROWS \}
FETCH \{ FIRST | NEXT \} \{ count \} \{ ROW | ROWS \} ONLY
\end{verbatim}

In this syntax, the \textit{start} or \textit{count} value is required by the standard to be a literal constant, a parameter, or a variable name; as a PostgreSQL extension, other expressions are allowed, but will generally need to be enclosed in parentheses to avoid ambiguity. If \textit{count} is omitted in a \texttt{FETCH} clause, it defaults to \texttt{1}. \texttt{ROW} and \texttt{ROWS} as well as \texttt{FIRST} and \texttt{NEXT} are noise words that don't influence the effects of these clauses. According to the standard, the \texttt{OFFSET} clause must come before the \texttt{FETCH} clause if both are present; but PostgreSQL is laxer and allows either order.

When using \texttt{LIMIT}, it is a good idea to use an \texttt{ORDER BY} clause that constrains the result rows into a unique order. Otherwise you will get an unpredictable subset of the query's rows — you might be asking for the tenth through twentieth rows, but tenth through twentieth in what ordering? You don't know what ordering unless you specify \texttt{ORDER BY}.

The query planner takes \texttt{LIMIT} into account when generating a query plan, so you are very likely to get different plans (yielding different row orders) depending on what you use for \texttt{LIMIT} and \texttt{OFFSET}. Thus, using different \texttt{LIMIT/OFFSET} values to select different subsets of a query result will give \textit{inconsistent results} unless you enforce a predictable result ordering with \texttt{ORDER BY}. This is not a bug; it is an inherent consequence of the fact that SQL does not promise to deliver the results of a query in any particular order unless \texttt{ORDER BY} is used to constrain the order.

It is even possible for repeated executions of the same \texttt{LIMIT} query to return different subsets of the rows of a table, if there is not an \texttt{ORDER BY} to enforce selection of a deterministic subset. Again, this is not a bug; determinism of the results is simply not guaranteed in such a case.

\section*{The Locking Clause}

\texttt{FOR UPDATE}, \texttt{FOR NO KEY UPDATE}, \texttt{FOR SHARE} and \texttt{FOR KEY SHARE} are locking clauses; they affect how \texttt{SELECT} locks rows as they are obtained from the table.

The locking clause has the general form

\begin{verbatim}
FOR lock_strength \[ OF table_name \[, \ldots \] \] \[ NOWAIT | SKIP LOCKED \]
\end{verbatim}
where \textit{lock\_strength} can be one of

\begin{itemize}
  \item \texttt{UPDATE}
  \item \texttt{NO KEY UPDATE}
  \item \texttt{SHARE}
  \item \texttt{KEY SHARE}
\end{itemize}

For more information on each row-level lock mode, refer to Section 13.3.2.

To prevent the operation from waiting for other transactions to commit, use either the \texttt{NOWAIT} or \texttt{SKIP LOCKED} option. With \texttt{NOWAIT}, the statement reports an error, rather than waiting, if a selected row cannot be locked immediately. With \texttt{SKIP LOCKED}, any selected rows that cannot be immediately locked are skipped. Skipping locked rows provides an inconsistent view of the data, so this is not suitable for general purpose work, but can be used to avoid lock contention with multiple consumers accessing a queue-like table. Note that \texttt{NOWAIT} and \texttt{SKIP LOCKED} apply only to the row-level lock(s) — the required \texttt{ROW SHARE} table-level lock is still taken in the ordinary way (see Chapter 13). You can use \texttt{LOCK} with the \texttt{NOWAIT} option first, if you need to acquire the table-level lock without waiting.

If specific tables are named in a locking clause, then only rows coming from those tables are locked; any other tables used in the \texttt{SELECT} are simply read as usual. A locking clause without a table list affects all tables used in the statement. If a locking clause is applied to a view or sub-query, it affects all tables used in the view or sub-query. However, these clauses do not apply to \texttt{WITH} queries referenced by the primary query. If you want row locking to occur within a \texttt{WITH} query, specify a locking clause within the \texttt{WITH} query.

Multiple locking clauses can be written if it is necessary to specify different locking behavior for different tables. If the same table is mentioned (or implicitly affected) by more than one locking clause, then it is processed as if it was only specified by the strongest one. Similarly, a table is processed as \texttt{NOWAIT} if that is specified in any of the clauses affecting it. Otherwise, it is processed as \texttt{SKIP LOCKED} if that is specified in any of the clauses affecting it.

The locking clauses cannot be used in contexts where returned rows cannot be clearly identified with individual table rows; for example they cannot be used with aggregation.

When a locking clause appears at the top level of a \texttt{SELECT} query, the rows that are locked are exactly those that are returned by the query; in the case of a join query, the rows locked are those that contribute to returned join rows. In addition, rows that satisfied the query conditions as of the query snapshot will be locked, although they will not be returned if they were updated after the snapshot and no longer satisfy the query conditions. If a \texttt{LIMIT} is used, locking stops once enough rows have been returned to satisfy the limit (but note that rows skipped over by \texttt{OFFSET} will get locked). Similarly, if a locking clause is used in a cursor's query, only rows actually fetched or stepped past by the cursor will be locked.

When a locking clause appears in a sub-\texttt{SELECT}, the rows locked are those returned to the outer query by the sub-query. This might involve fewer rows than inspection of the sub-query alone would suggest, since conditions from the outer query might be used to optimize execution of the sub-query. For example,

\begin{verbatim}
SELECT * FROM (SELECT * FROM mytable FOR UPDATE) ss WHERE col1 = 5;
\end{verbatim}

will lock only rows having \texttt{col1 = 5}, even though that condition is not textually within the sub-query.

Previous releases failed to preserve a lock which is upgraded by a later savepoint. For example, this code:

\begin{verbatim}
SELECT * FROM (SELECT * FROM mytable FOR UPDATE) ss WHERE col1 = 5;
\end{verbatim}

will lock only rows having \texttt{col1 = 5}, even though that condition is not textually within the sub-query.
BEGIN;
SELECT * FROM mytable WHERE key = 1 FOR UPDATE;
SAVEPOINT s;
UPDATE mytable SET ... WHERE key = 1;
ROLLBACK TO s;

would fail to preserve the FOR UPDATE lock after the ROLLBACK TO. This has been fixed in release 9.3.

### Caution

It is possible for a SELECT command running at the READ COMMITTED transaction isolation level and using ORDER BY and a locking clause to return rows out of order. This is because ORDER BY is applied first. The command sorts the result, but might then block trying to obtain a lock on one or more of the rows. Once the SELECT unblocks, some of the ordering column values might have been modified, leading to those rows appearing to be out of order (though they are in order in terms of the original column values). This can be worked around at need by placing the FOR UPDATE/SHARE clause in a sub-query, for example

```sql
SELECT * FROM (SELECT * FROM mytable FOR UPDATE) ss ORDER BY column1;
```

Note that this will result in locking all rows of mytable, whereas FOR UPDATE at the top level would lock only the actually returned rows. This can make for a significant performance difference, particularly if the ORDER BY is combined with LIMIT or other restrictions. So this technique is recommended only if concurrent updates of the ordering columns are expected and a strictly sorted result is required.

At the REPEATABLE READ or SERIALIZABLE transaction isolation level this would cause a serialization failure (with a SQLSTATE of '40001'), so there is no possibility of receiving rows out of order under these isolation levels.

### TABLE Command

The command

```
TABLE name
```

is equivalent to

```
SELECT * FROM name
```

It can be used as a top-level command or as a space-saving syntax variant in parts of complex queries. Only the WITH, UNION, INTERSECT, EXCEPT, ORDER BY, LIMIT, OFFSET, FETCH and FOR locking clauses can be used with TABLE; the WHERE clause and any form of aggregation cannot be used.

### Examples

To join the table films with the table distributors:
SELECT f.title, f.did, d.name, f.date_prod, f.kind
FROM distributors d, films f
WHERE f.did = d.did

<table>
<thead>
<tr>
<th>title</th>
<th>did</th>
<th>name</th>
<th>date_prod</th>
<th>kind</th>
</tr>
</thead>
<tbody>
<tr>
<td>The Third Man</td>
<td>101</td>
<td>British Lion</td>
<td>1949-12-23</td>
<td>Drama</td>
</tr>
<tr>
<td>The African Queen</td>
<td>101</td>
<td>British Lion</td>
<td>1951-08-11</td>
<td>Romantic</td>
</tr>
</tbody>
</table>

To sum the column `len` of all films and group the results by `kind`:

```sql
SELECT kind, sum(len) AS total
FROM films
GROUP BY kind;
```

<table>
<thead>
<tr>
<th>kind</th>
<th>total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Action</td>
<td>07:34</td>
</tr>
<tr>
<td>Comedy</td>
<td>02:58</td>
</tr>
<tr>
<td>Drama</td>
<td>14:28</td>
</tr>
<tr>
<td>Musical</td>
<td>06:42</td>
</tr>
<tr>
<td>Romantic</td>
<td>04:38</td>
</tr>
</tbody>
</table>

To sum the column `len` of all films, group the results by `kind` and show those group totals that are less than 5 hours:

```sql
SELECT kind, sum(len) AS total
FROM films
GROUP BY kind
HAVING sum(len) < interval '5 hours';
```

<table>
<thead>
<tr>
<th>kind</th>
<th>total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Comedy</td>
<td>02:58</td>
</tr>
<tr>
<td>Romantic</td>
<td>04:38</td>
</tr>
</tbody>
</table>

The following two examples are identical ways of sorting the individual results according to the contents of the second column (`name`):

```sql
SELECT * FROM distributors ORDER BY name;
SELECT * FROM distributors ORDER BY 2;
```

<table>
<thead>
<tr>
<th>did</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>109</td>
<td>20th Century Fox</td>
</tr>
<tr>
<td>110</td>
<td>Bavaria Atelier</td>
</tr>
<tr>
<td>101</td>
<td>British Lion</td>
</tr>
<tr>
<td>107</td>
<td>Columbia</td>
</tr>
<tr>
<td>102</td>
<td>Jean Luc Godard</td>
</tr>
<tr>
<td>113</td>
<td>Luso films</td>
</tr>
<tr>
<td>104</td>
<td>Mosfilm</td>
</tr>
<tr>
<td>103</td>
<td>Paramount</td>
</tr>
<tr>
<td>106</td>
<td>Toho</td>
</tr>
<tr>
<td>105</td>
<td>United Artists</td>
</tr>
<tr>
<td>111</td>
<td>Walt Disney</td>
</tr>
<tr>
<td>112</td>
<td>Warner Bros.</td>
</tr>
<tr>
<td>108</td>
<td>Westward</td>
</tr>
</tbody>
</table>
The next example shows how to obtain the union of the tables distributors and actors, restricting the results to those that begin with the letter W in each table. Only distinct rows are wanted, so the key word ALL is omitted.

```sql
SELECT distributors.name
FROM distributors
WHERE distributors.name LIKE 'W%'
UNION
SELECT actors.name
FROM actors
WHERE actors.name LIKE 'W%';
```

This example shows how to use a function in the FROM clause, both with and without a column definition list:

```sql
CREATE FUNCTION distributors(int) RETURNS SETOF distributors AS $$
SELECT * FROM distributors WHERE did = $1;
$$ LANGUAGE SQL;

SELECT * FROM distributors(111);
  did |    name
-----+------------
 111 | Walt Disney

CREATE FUNCTION distributors_2(int) RETURNS SETOF record AS $$
SELECT * FROM distributors WHERE did = $1;
$$ LANGUAGE SQL;

SELECT * FROM distributors_2(111) AS (f1 int, f2 text);
  f1  |     f2
------|------------
 111  | Walt Disney
```

Here is an example of a function with an ordinality column added:

```sql
SELECT * FROM unnest(ARRAY['a','b','c','d','e','f']) WITH ORDINALITY;
   unnest | ordinality
---------|----------
"a"      |    1
"b"      |    2
"c"      |    3
"d"      |    4
"e"      |    5
"f"      |    6
```
This example shows how to use a simple WITH clause:

```sql
WITH t AS (
    SELECT random() as x FROM generate_series(1, 3)
) 
SELECT * FROM t 
UNION ALL 
SELECT * FROM t ;
```

Notice that the WITH query was evaluated only once, so that we got two sets of the same three random values.

This example uses WITH RECURSIVE to find all subordinates (direct or indirect) of the employee Mary, and their level of indirectness, from a table that shows only direct subordinates:

```sql
WITH RECURSIVE employee_recursive(distance, employee_name, manager_name) AS ( 
    SELECT 1, employee_name, manager_name 
    FROM employee 
    WHERE manager_name = 'Mary' 
    UNION ALL 
    SELECT er.distance + 1, e.employee_name, e.manager_name 
    FROM employee_recursive er, employee e 
    WHERE er.employee_name = e.manager_name 
) 
SELECT distance, employee_name FROM employee_recursive;
```

Notice the typical form of recursive queries: an initial condition, followed by UNION, followed by the recursive part of the query. Be sure that the recursive part of the query will eventually return no tuples, or else the query will loop indefinitely. (See Section 7.8 for more examples.)

This example uses LATERAL to apply a set-returning function get_product_names() for each row of the manufacturers table:

```sql
SELECT m.name AS mname, pname 
FROM manufacturers m, LATERAL get_product_names(m.id) pname;
```

Manufacturers not currently having any products would not appear in the result, since it is an inner join. If we wished to include the names of such manufacturers in the result, we could do:
SELECT m.name AS mname, pname 
FROM manufacturers m LEFT JOIN LATERAL get_product_names(m.id) 
pname ON true;

Compatibility

Of course, the SELECT statement is compatible with the SQL standard. But there are some extensions and some missing features.

Omitted FROM Clauses

PostgreSQL allows one to omit the FROM clause. It has a straightforward use to compute the results of simple expressions:

```
SELECT 2+2;
```

<table>
<thead>
<tr>
<th>?column?</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
</tr>
</tbody>
</table>

Some other SQL databases cannot do this except by introducing a dummy one-row table from which to do the SELECT.

Note that if a FROM clause is not specified, the query cannot reference any database tables. For example, the following query is invalid:

```
SELECT distributors.* WHERE distributors.name = 'Westward';
```

PostgreSQL releases prior to 8.1 would accept queries of this form, and add an implicit entry to the query’s FROM clause for each table referenced by the query. This is no longer allowed.

Empty SELECT Lists

The list of output expressions after SELECT can be empty, producing a zero-column result table. This is not valid syntax according to the SQL standard. PostgreSQL allows it to be consistent with allowing zero-column tables. However, an empty list is not allowed when DISTINCT is used.

Omitting the AS Key Word

In the SQL standard, the optional key word AS can be omitted before an output column name whenever the new column name is a valid column name (that is, not the same as any reserved keyword). PostgreSQL is slightly more restrictive: AS is required if the new column name matches any keyword at all, reserved or not. Recommended practice is to use AS or double-quote output column names, to prevent any possible conflict against future keyword additions.

In FROM items, both the standard and PostgreSQL allow AS to be omitted before an alias that is an unreserved keyword. But this is impractical for output column names, because of syntactic ambiguities.

ONLY and Inheritance

The SQL standard requires parentheses around the table name when writing ONLY, for example

```
SELECT * FROM ONLY (tab1), ONLY (tab2) WHERE ... . . . PostgreSQL considers these parentheses to be optional.
```

PostgreSQL allows a trailing * to be written to explicitly specify the non-ONLY behavior of including child tables. The standard does not allow this.
(These points apply equally to all SQL commands supporting the ONLY option.)

**TABLESAMPLE Clause Restrictions**

The TABLESAMPLE clause is currently accepted only on regular tables and materialized views. According to the SQL standard it should be possible to apply it to any FROM item.

**Function Calls in FROM**

PostgreSQL allows a function call to be written directly as a member of the FROM list. In the SQL standard it would be necessary to wrap such a function call in a sub-SELECT; that is, the syntax FROM \texttt{func(...) alias} is approximately equivalent to FROM LATERAL (SELECT func(...)) alias. Note that LATERAL is considered to be implicit; this is because the standard requires LATERAL semantics for an UNNEST() item in FROM. PostgreSQL treats UNNEST() the same as other set-returning functions.

**Namespace Available to GROUP BY and ORDER BY**

In the SQL-92 standard, an ORDER BY clause can only use output column names or numbers, while a GROUP BY clause can only use expressions based on input column names. PostgreSQL extends each of these clauses to allow the other choice as well (but it uses the standard's interpretation if there is ambiguity). PostgreSQL also allows both clauses to specify arbitrary expressions. Note that names appearing in an expression will always be taken as input-column names, not as output-column names.

SQL:1999 and later use a slightly different definition which is not entirely upward compatible with SQL-92. In most cases, however, PostgreSQL will interpret an ORDER BY or GROUP BY expression the same way SQL:1999 does.

**Functional Dependencies**

PostgreSQL recognizes functional dependency (allowing columns to be omitted from GROUP BY) only when a table's primary key is included in the GROUP BY list. The SQL standard specifies additional conditions that should be recognized.

**LIMIT and OFFSET**

The clauses LIMIT and OFFSET are PostgreSQL-specific syntax, also used by MySQL. The SQL:2008 standard has introduced the clauses OFFSET ... FETCH \{FIRST|NEXT\} ... for the same functionality, as shown above in LIMIT Clause. This syntax is also used by IBM DB2. (Applications written for Oracle frequently use a workaround involving the automatically generated rownum column, which is not available in PostgreSQL, to implement the effects of these clauses.)

**FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, FOR KEY SHARE**

Although FOR UPDATE appears in the SQL standard, the standard allows it only as an option of DECLARE CURSOR. PostgreSQL allows it in any SELECT query as well as in sub-SELECTs, but this is an extension. The FOR NO KEY UPDATE, FOR SHARE and FOR KEY SHARE variants, as well as the NOWAIT and SKIP LOCKED options, do not appear in the standard.

**Data-Modifying Statements in WITH**

PostgreSQL allows INSERT, UPDATE, and DELETE to be used as WITH queries. This is not found in the SQL standard.

**Nonstandard Clauses**

\texttt{DISTINCT ON ( ... )} is an extension of the SQL standard.
SELECT

ROWS FROM ( ... ) is an extension of the SQL standard.
SELECT INTO

SELECT INTO — define a new table from the results of a query

Synopsis

\[
\begin{align*}
\text{[ WITH [ RECURSIVE ] with_query [, ...] ]} \\
\text{SELECT [ ALL | DISTINCT [ ON ( expression [, ...] ) ] ]} \\
\quad [ \text{ expression [ [ AS ] output_name ] [, ...]} ] \\
\quad \text{INTO [ TEMPORARY | TEMP | UNLOGGED ] [ TABLE ] new_table} \\
\quad [ \text{ FROM from_item [, ...] } ] \\
\quad [ \text{ WHERE condition } ] \\
\quad [ \text{ GROUP BY expression [, ...] } ] \\
\quad [ \text{ HAVING condition [, ...] } ] \\
\quad [ \text{ WINDOW window_name AS ( window_definition ) [, ...] } ] \\
\quad [ \{ \text{ UNION | INTERSECT | EXCEPT } \} [ \text{ ALL | DISTINCT } \] \text{ select } ] \\
\quad [ \text{ ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS} \\
\quad \text{( FIRST | LAST ) ] [, ...] ] \\
\quad [ \text{ LIMIT { count | ALL } ] } \\
\quad [ \text{ OFFSET start [ ROW | ROWS ] } ] \\
\quad [ \text{ FETCH { FIRST | NEXT } [ count ] [ ROW | ROWS ] ONLY } ] \\
\quad [ \text{ FOR { UPDATE | SHARE } [ OF table_name [, ...] ] [ NOWAIT } ] \\
\quad [\ldots] ]
\end{align*}
\]

Description

SELECT INTO creates a new table and fills it with data computed by a query. The data is not returned to the client, as it is with a normal SELECT. The new table’s columns have the names and data types associated with the output columns of the SELECT.

Parameters

TEMPORARY or TEMP

If specified, the table is created as a temporary table. Refer to CREATE TABLE for details.

UNLOGGED

If specified, the table is created as an unlogged table. Refer to CREATE TABLE for details.

new_table

The name (optionally schema-qualified) of the table to be created.

All other parameters are described in detail under SELECT.

Notes

CREATE TABLE AS is functionally similar to SELECT INTO. CREATE TABLE AS is the recommended syntax, since this form of SELECT INTO is not available in ECPG or PL/pgSQL, because they interpret the INTO clause differently. Furthermore, CREATE TABLE AS offers a superset of the functionality provided by SELECT INTO.

To add OIDs to the table created by SELECT INTO, enable the default with oids configuration variable. Alternatively, CREATE TABLE AS can be used with the WITH OIDS clause.
Examples

Create a new table `films_recent` consisting of only recent entries from the table `films`:

```sql
SELECT * INTO films_recent FROM films WHERE date_prod >= '2002-01-01';
```

Compatibility

The SQL standard uses `SELECT INTO` to represent selecting values into scalar variables of a host program, rather than creating a new table. This indeed is the usage found in ECPG (see Chapter 36) and PL/pgSQL (see Chapter 43). The PostgreSQL usage of `SELECT INTO` to represent table creation is historical. It is best to use `CREATE TABLE AS` for this purpose in new code.

See Also

CREATE TABLE AS
SET

SET — change a run-time parameter

Synopsis

SET [ SESSION | LOCAL ] configuration_parameter { TO | = } { value | 'value' | DEFAULT }
SET [ SESSION | LOCAL ] TIME ZONE { timezone | LOCAL | DEFAULT }

Description

The SET command changes run-time configuration parameters. Many of the run-time parameters listed in Chapter 19 can be changed on-the-fly with SET. (But some require superuser privileges to change, and others cannot be changed after server or session start.) SET only affects the value used by the current session.

If SET (or equivalently SET SESSION) is issued within a transaction that is later aborted, the effects of the SET command disappear when the transaction is rolled back. Once the surrounding transaction is committed, the effects will persist until the end of the session, unless overridden by another SET.

The effects of SET LOCAL last only till the end of the current transaction, whether committed or not. A special case is SET followed by SET LOCAL within a single transaction: the SET LOCAL value will be seen until the end of the transaction, but afterwards (if the transaction is committed) the SET value will take effect.

The effects of SET or SET LOCAL are also canceled by rolling back to a savepoint that is earlier than the command.

If SET LOCAL is used within a function that has a SET option for the same variable (see CREATE FUNCTION), the effects of the SET LOCAL command disappear at function exit; that is, the value in effect when the function was called is restored anyway. This allows SET LOCAL to be used for dynamic or repeated changes of a parameter within a function, while still having the convenience of using the SET option to save and restore the caller's value. However, a regular SET command overrides any surrounding function's SET option; its effects will persist unless rolled back.

Note

In PostgreSQL versions 8.0 through 8.2, the effects of a SET LOCAL would be canceled by releasing an earlier savepoint, or by successful exit from a PL/pgSQL exception block. This behavior has been changed because it was deemed unintuitive.

Parameters

SESSION

Specifies that the command takes effect for the current session. (This is the default if neither SESSION nor LOCAL appears.)

LOCAL

Specifies that the command takes effect for only the current transaction. After COMMIT or ROLLBACK, the session-level setting takes effect again. Issuing this outside of a transaction block emits a warning and otherwise has no effect.
configuration_parameter

Name of a settable run-time parameter. Available parameters are documented in Chapter 19 and below.

value

New value of parameter. Values can be specified as string constants, identifiers, numbers, or comma-separated lists of these, as appropriate for the particular parameter. DEFAULT can be written to specify resetting the parameter to its default value (that is, whatever value it would have had if no SET had been executed in the current session).

Besides the configuration parameters documented in Chapter 19, there are a few that can only be adjusted using the SET command or that have a special syntax:

SCHEMA

SET SCHEMA 'value' is an alias for SET search_path TO value. Only one schema can be specified using this syntax.

NAMES

SET NAMES value is an alias for SET client_encoding TO value.

SEED

Sets the internal seed for the random number generator (the function random). Allowed values are floating-point numbers between -1 and 1, which are then multiplied by $2^{31}$-1.

The seed can also be set by invoking the function setseed:

SELECT setseed(value);

TIME ZONE

SET TIME ZONE value is an alias for SET timezone TO value. The syntax SET TIME ZONE allows special syntax for the time zone specification. Here are examples of valid values:

'PST8PDT'

The time zone for Berkeley, California.

'Europe/Rome'

The time zone for Italy.

-7

The time zone 7 hours west from UTC (equivalent to PDT). Positive values are east from UTC.

INTERVAL '-08:00' HOUR TO MINUTE

The time zone 8 hours west from UTC (equivalent to PST).

LOCAL DEFAULT

Set the time zone to your local time zone (that is, the server’s default value of timezone).

Timezone settings given as numbers or intervals are internally translated to POSIX timezone syntax. For example, after SET TIME ZONE -7, SHOW TIME ZONE would report <-07>+07.
SET

See Section 8.5.3 for more information about time zones.

Notes

The function `set_config` provides equivalent functionality; see Section 9.26. Also, it is possible to UPDATE the `pg_settings` system view to perform the equivalent of `SET`.

Examples

Set the schema search path:

```
SET search_path TO my_schema, public;
```

Set the style of date to traditional POSTGRES with “day before month” input convention:

```
SET datestyle TO postgres, dmy;
```

Set the time zone for Berkeley, California:

```
SET TIME ZONE 'PST8PDT';
```

Set the time zone for Italy:

```
SET TIME ZONE 'Europe/Rome';
```

Compatibility

`SET TIME ZONE` extends syntax defined in the SQL standard. The standard allows only numeric time zone offsets while PostgreSQL allows more flexible time-zone specifications. All other `SET` features are PostgreSQL extensions.

See Also

`RESET`, `SHOW`
SET CONSTRAINTS

SET CONSTRAINTS — set constraint check timing for the current transaction

Synopsis

SET CONSTRAINTS { ALL | name [, ...] } { DEFERRED | IMMEDIATE }

Description

SET CONSTRAINTS sets the behavior of constraint checking within the current transaction. IMMEDIATE constraints are checked at the end of each statement. DEFERRED constraints are not checked until transaction commit. Each constraint has its own IMMEDIATE or DEFERRED mode.

Upon creation, a constraint is given one of three characteristics: DEFERRABLE INITIALLY DEFERRED, DEFERRABLE INITIALLY IMMEDIATE, or NOT DEFERRABLE. The third class is always IMMEDIATE and is not affected by the SET CONSTRAINTS command. The first two classes start every transaction in the indicated mode, but their behavior can be changed within a transaction by SET CONSTRAINTS.

SET CONSTRAINTS with a list of constraint names changes the mode of just those constraints (which must all be deferrable). Each constraint name can be schema-qualified. The current schema search path is used to find the first matching name if no schema name is specified. SET CONSTRAINTS ALL changes the mode of all deferrable constraints.

When SET CONSTRAINTS changes the mode of a constraint from DEFERRED to IMMEDIATE, the new mode takes effect retroactively: any outstanding data modifications that would have been checked at the end of the transaction are instead checked during the execution of the SET CONSTRAINTS command. If any such constraint is violated, the SET CONSTRAINTS fails (and does not change the constraint mode). Thus, SET CONSTRAINTS can be used to force checking of constraints to occur at a specific point in a transaction.

Currently, only UNIQUE, PRIMARY KEY, REFERENCES (foreign key), and EXCLUDE constraints are affected by this setting. NOT NULL and CHECK constraints are always checked immediately when a row is inserted or modified (not at the end of the statement). Uniqueness and exclusion constraints that have not been declared DEFERRABLE are also checked immediately.

The firing of triggers that are declared as “constraint triggers” is also controlled by this setting — they fire at the same time that the associated constraint should be checked.

Notes

Because PostgreSQL does not require constraint names to be unique within a schema (but only per-table), it is possible that there is more than one match for a specified constraint name. In this case SET CONSTRAINTS will act on all matches. For a non-schema-qualified name, once a match or matches have been found in some schema in the search path, schemas appearing later in the path are not searched.

This command only alters the behavior of constraints within the current transaction. Issuing this outside of a transaction block emits a warning and otherwise has no effect.

Compatibility

This command complies with the behavior defined in the SQL standard, except for the limitation that, in PostgreSQL, it does not apply to NOT NULL and CHECK constraints. Also, PostgreSQL checks non-deferrable uniqueness constraints immediately, not at end of statement as the standard would suggest.
**SET ROLE**

SET ROLE — set the current user identifier of the current session

**Synopsis**

```plaintext
SET [ SESSION | LOCAL ] ROLE role_name
SET [ SESSION | LOCAL ] ROLE NONE
RESET ROLE
```

**Description**

This command sets the current user identifier of the current SQL session to be `role_name`. The role name can be written as either an identifier or a string literal. After `SET ROLE`, permissions checking for SQL commands is carried out as though the named role were the one that had logged in originally.

The specified `role_name` must be a role that the current session user is a member of. (If the session user is a superuser, any role can be selected.)

The `SESSION` and `LOCAL` modifiers act the same as for the regular `SET` command.

The `NONE` and `RESET` forms reset the current user identifier to be the current session user identifier. These forms can be executed by any user.

**Notes**

Using this command, it is possible to either add privileges or restrict one's privileges. If the session user role has the `INHERITS` attribute, then it automatically has all the privileges of every role that it could `SET ROLE` to; in this case `SET ROLE` effectively drops all the privileges assigned directly to the session user and to the other roles it is a member of, leaving only the privileges available to the named role. On the other hand, if the session user role has the `NOINHERITS` attribute, `SET ROLE` drops the privileges assigned directly to the session user and instead acquires the privileges available to the named role.

In particular, when a superuser chooses to `SET ROLE` to a non-superuser role, they lose their superuser privileges.

`SET ROLE` has effects comparable to `SET SESSION AUTHORIZATION`, but the privilege checks involved are quite different. Also, `SET SESSION AUTHORIZATION` determines which roles are allowable for later `SET ROLE` commands, whereas changing roles with `SET ROLE` does not change the set of roles allowed to a later `SET ROLE`.

`SET ROLE` does not process session variables as specified by the role's `ALTER ROLE` settings; this only happens during login.

`SET ROLE` cannot be used within a `SECURITY DEFINER` function.

**Examples**

```sql
SELECT SESSION_USER, CURRENT_USER;
```

```
session_user | current_user
--------------+--------------
1791
```
peter | peter
SET ROLE 'paul';

SELECT SESSION_USER, CURRENT_USER;

<table>
<thead>
<tr>
<th>session_user</th>
<th>current_user</th>
</tr>
</thead>
<tbody>
<tr>
<td>peter</td>
<td>paul</td>
</tr>
</tbody>
</table>

**Compatibility**

PostgreSQL allows identifier syntax ("rolename"), while the SQL standard requires the role name to be written as a string literal. SQL does not allow this command during a transaction; PostgreSQL does not make this restriction because there is no reason to. The `SESSION` and `LOCAL` modifiers are a PostgreSQL extension, as is the `RESET` syntax.

**See Also**

SET SESSION AUTHORIZATION
SET SESSION AUTHORIZATION

SET SESSION AUTHORIZATION — set the session user identifier and the current user identifier of the current session

Synopsis

```
SET [ SESSION | LOCAL ] SESSION AUTHORIZATION user_name
SET [ SESSION | LOCAL ] SESSION AUTHORIZATION DEFAULT
RESET SESSION AUTHORIZATION
```

Description

This command sets the session user identifier and the current user identifier of the current SQL session to be `user_name`. The user name can be written as either an identifier or a string literal. Using this command, it is possible, for example, to temporarily become an unprivileged user and later switch back to being a superuser.

The session user identifier is initially set to the (possibly authenticated) user name provided by the client. The current user identifier is normally equal to the session user identifier, but might change temporarily in the context of SECURITY DEFINER functions and similar mechanisms; it can also be changed by SET ROLE. The current user identifier is relevant for permission checking.

The session user identifier can be changed only if the initial session user (the authenticated user) had the superuser privilege. Otherwise, the command is accepted only if it specifies the authenticated user name.

The `SESSION` and `LOCAL` modifiers act the same as for the regular SET command.

The `DEFAULT` and `RESET` forms reset the session and current user identifiers to be the originally authenticated user name. These forms can be executed by any user.

Notes

`SET SESSION AUTHORIZATION` cannot be used within a SECURITY DEFINER function.

Examples

```
SELECT SESSION_USER, CURRENT_USER;

+-----------+-----------+
| session_user | current_user |
+-----------+-----------+
| peter     | peter     |

SET SESSION AUTHORIZATION 'paul';

SELECT SESSION_USER, CURRENT_USER;

+-----------+-----------+
| session_user | current_user |
+-----------+-----------+
| paul      | paul      |
```
Compatibility

The SQL standard allows some other expressions to appear in place of the literal \texttt{user\_name}, but these options are not important in practice. PostgreSQL allows identifier syntax ("\texttt{username}"), which SQL does not. SQL does not allow this command during a transaction; PostgreSQL does not make this restriction because there is no reason to. The \texttt{SESSION} and \texttt{LOCAL} modifiers are a PostgreSQL extension, as is the \texttt{RESET} syntax.

The privileges necessary to execute this command are left implementation-defined by the standard.

See Also

\texttt{SET ROLE}
**SET TRANSACTION**

SET TRANSACTION — set the characteristics of the current transaction

**Synopsis**

```sql
SET TRANSACTION transaction_mode [, ...]
SET TRANSACTION SNAPSHOT snapshot_id
SET SESSION CHARACTERISTICS AS TRANSACTION transaction_mode [, ...]
```

where `transaction_mode` is one of:

- ISOLATION LEVEL { SERIALizable | REPEATable READ | READ COMMITTED | READ UNCOMMITTED }
- READ WRITE | READ ONLY
- [ NOT ] DEFERRABLE

**Description**

The `SET TRANSACTION` command sets the characteristics of the current transaction. It has no effect on any subsequent transactions. `SET SESSION CHARACTERISTICS` sets the default transaction characteristics for subsequent transactions of a session. These defaults can be overridden by `SET TRANSACTION` for an individual transaction.

The available transaction characteristics are the transaction isolation level, the transaction access mode (read/write or read-only), and the deferrable mode. In addition, a snapshot can be selected, though only for the current transaction, not as a session default.

The isolation level of a transaction determines what data the transaction can see when other transactions are running concurrently:

**READ COMMITTED**

A statement can only see rows committed before it began. This is the default.

**REPEATABLE READ**

All statements of the current transaction can only see rows committed before the first query or data-modification statement was executed in this transaction.

**SERIALIZABLE**

All statements of the current transaction can only see rows committed before the first query or data-modification statement was executed in this transaction. If a pattern of reads and writes among concurrent serializable transactions would create a situation which could not have occurred for any serial (one-at-a-time) execution of those transactions, one of them will be rolled back with a `serialization_failure` error.

The SQL standard defines one additional level, **READ UNCOMMITTED**. In PostgreSQL READ UNCOMMITTED is treated as **READ COMMITTED**.

The transaction isolation level cannot be changed after the first query or data-modification statement (SELECT, INSERT, DELETE, UPDATE, FETCH, or COPY) of a transaction has been executed. See Chapter 13 for more information about transaction isolation and concurrency control.

The transaction access mode determines whether the transaction is read/write or read-only. Read/write is the default. When a transaction is read-only, the following SQL commands are disallowed: INSERT,
SET TRANSACTION

UPDATE, DELETE, and COPY FROM if the table they would write to is not a temporary table; all CREATE, ALTER, and DROP commands; COMMENT, GRANT, REVOKE, TRUNCATE; and EXPLAIN ANALYZE and EXECUTE if the command they would execute is among those listed. This is a high-level notion of read-only that does not prevent all writes to disk.

The DEFERRABLE transaction property has no effect unless the transaction is also SERIALIZABLE and READ ONLY. When all three of these properties are selected for a transaction, the transaction may block when first acquiring its snapshot, after which it is able to run without the normal overhead of a SERIALIZABLE transaction and without any risk of contributing to or being canceled by a serialization failure. This mode is well suited for long-running reports or backups.

The SET TRANSACTION SNAPSHOT command allows a new transaction to run with the same snapshot as an existing transaction. The pre-existing transaction must have exported its snapshot with the pg_export_snapshot function (see Section 9.26.5). That function returns a snapshot identifier, which must be given to SET TRANSACTION SNAPSHOT to specify which snapshot is to be imported. The identifier must be written as a string literal in this command, for example '0000003A1-1'. SET TRANSACTION SNAPSHOT can only be executed at the start of a transaction, before the first query or data-modification statement (SELECT, INSERT, DELETE, UPDATE, FETCH, or COPY) of the transaction. Furthermore, the transaction must already be set to SERIALIZABLE or REPEATABLE READ isolation level (otherwise, the snapshot would be discarded immediately, since READ COMMITTED mode takes a new snapshot for each command). If the importing transaction uses SERIALIZABLE isolation level, then the transaction that exported the snapshot must also use that isolation level. Also, a non-read-only serializable transaction cannot import a snapshot from a read-only transaction.

Notes

If SET TRANSACTION is executed without a prior START TRANSACTION or BEGIN, it emits a warning and otherwise has no effect.

It is possible to dispense with SET TRANSACTION by instead specifying the desired transaction_modes in BEGIN or START TRANSACTION. But that option is not available for SET TRANSACTION SNAPSHOT.

The session default transaction modes can also be set by setting the configuration parameters default_transaction_isolation, default_transaction_read_only, and default_transaction_deferrable. (In fact SET SESSION CHARACTERISTICS is just a verbose equivalent for setting these variables with SET.) This means the defaults can be set in the configuration file, via ALTER DATABASE, etc. Consult Chapter 19 for more information.

Examples

To begin a new transaction with the same snapshot as an already existing transaction, first export the snapshot from the existing transaction. That will return the snapshot identifier, for example:

BEGIN TRANSACTION ISOLATION LEVEL REPEATABLE READ;
SELECT pg_export_snapshot();
pg_export_snapshot
-------------------
00000003-0000001B-1
(1 row)

Then give the snapshot identifier in a SET TRANSACTION SNAPSHOT command at the beginning of the newly opened transaction:

BEGIN TRANSACTION ISOLATION LEVEL REPEATABLE READ;
SET TRANSACTION SNAPSHOT '00000003-0000001B-1';
Compatibility

These commands are defined in the SQL standard, except for the DEFERRABLE transaction mode and the SET TRANSACTION SNAPSHOT form, which are PostgreSQL extensions.

SERIALIZABLE is the default transaction isolation level in the standard. In PostgreSQL the default is ordinarily READ COMMITTED, but you can change it as mentioned above.

In the SQL standard, there is one other transaction characteristic that can be set with these commands: the size of the diagnostics area. This concept is specific to embedded SQL, and therefore is not implemented in the PostgreSQL server.

The SQL standard requires commas between successive transaction_modes, but for historical reasons PostgreSQL allows the commas to be omitted.
SHOW

SHOW — show the value of a run-time parameter

Synopsis

SHOW name
SHOW ALL

Description

SHOW will display the current setting of run-time parameters. These variables can be set using the SET statement, by editing the postgresql.conf configuration file, through the PGOPTIONS environmental variable (when using libpq or a libpq-based application), or through command-line flags when starting the postgres server. See Chapter 19 for details.

Parameters

name

The name of a run-time parameter. Available parameters are documented in Chapter 19 and on the SET reference page. In addition, there are a few parameters that can be shown but not set:

SERVER_VERSION

Shows the server's version number.

SERVER_ENCODING

Shows the server-side character set encoding. At present, this parameter can be shown but not set, because the encoding is determined at database creation time.

LC_COLLATE

Shows the database's locale setting for collation (text ordering). At present, this parameter can be shown but not set, because the setting is determined at database creation time.

LC_CTYPE

Shows the database's locale setting for character classification. At present, this parameter can be shown but not set, because the setting is determined at database creation time.

IS_SUPERUSER

True if the current role has superuser privileges.

ALL

Show the values of all configuration parameters, with descriptions.

Notes

The function current_setting produces equivalent output; see Section 9.26. Also, the pg_settings system view produces the same information.
Examples

Show the current setting of the parameter `DateStyle`:

```
SHOW DateStyle;
DateStyle
----------
ISO, MDY
(1 row)
```

Show the current setting of the parameter `geqo`:

```
SHOW geqo;
geqo
------
on
(1 row)
```

Show all settings:

```
SHOW ALL;
name         | setting | description
-------------------------+---------|
+--------------------------------------------------
allow_system_table_mods | off     | Allows modifications of the structure of ... . . .
xmloption               | content | Sets whether XML data in implicit parsing ... zero_damaged_pages | off | Continues processing past damaged page headers.
(196 rows)
```

Compatibility

The `SHOW` command is a PostgreSQL extension.

See Also

`SET`, `RESET`
START TRANSACTION

START TRANSACTION — start a transaction block

Synopsis

START TRANSACTION [ transaction_mode [, ... ] ]

where transaction_mode is one of:

- ISOLATION LEVEL { SERIALIZABLE | REPEATABLE READ | READ COMMITTED | READ UNCOMMITTED }
- READ WRITE | READ ONLY
- [ NOT ] DEFERRABLE

Description

This command begins a new transaction block. If the isolation level, read/write mode, or deferrable mode is specified, the new transaction has those characteristics, as if SET TRANSACTION was executed. This is the same as the BEGIN command.

Parameters

Refer to SET TRANSACTION for information on the meaning of the parameters to this statement.

Compatibility

In the standard, it is not necessary to issue START TRANSACTION to start a transaction block: any SQL command implicitly begins a block. PostgreSQL’s behavior can be seen as implicitly issuing a COMMIT after each command that does not follow START TRANSACTION (or BEGIN), and it is therefore often called “autocommit”. Other relational database systems might offer an autocommit feature as a convenience.

The DEFERRABLE transaction_mode is a PostgreSQL language extension.

The SQL standard requires commas between successive transaction_modes, but for historical reasons PostgreSQL allows the commas to be omitted.

See also the compatibility section of SET TRANSACTION.

See Also

BEGIN, COMMIT, ROLLBACK, SAVEPOINT, SET TRANSACTION
TRUNCATE

TRUNCATE — empty a table or set of tables

Synopsis

TRUNCATE [ TABLE ] [ ONLY ] name [ * ] [, ... ]
[ RESTART IDENTITY | CONTINUE IDENTITY ] [ CASCADE | RESTRICT ]

Description

TRUNCATE quickly removes all rows from a set of tables. It has the same effect as an unqualified DELETE on each table, but since it does not actually scan the tables it is faster. Furthermore, it reclaims disk space immediately, rather than requiring a subsequent VACUUM operation. This is most useful on large tables.

Parameters

name

The name (optionally schema-qualified) of a table to truncate. If ONLY is specified before the table name, only that table is truncated. If ONLY is not specified, the table and all its descendant tables (if any) are truncated. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

RESTART IDENTITY

Automatically restart sequences owned by columns of the truncated table(s).

CONTINUE IDENTITY

Do not change the values of sequences. This is the default.

CASCADE

Automatically truncate all tables that have foreign-key references to any of the named tables, or to any tables added to the group due to CASCADE.

RESTRICT

Refuse to truncate if any of the tables have foreign-key references from tables that are not listed in the command. This is the default.

Notes

You must have the TRUNCATE privilege on a table to truncate it.

TRUNCATE acquires an ACCESS EXCLUSIVE lock on each table it operates on, which blocks all other concurrent operations on the table. When RESTART IDENTITY is specified, any sequences that are to be restarted are likewise locked exclusively. If concurrent access to a table is required, then the DELETE command should be used instead.

TRUNCATE cannot be used on a table that has foreign-key references from other tables, unless all such tables are also truncated in the same command. Checking validity in such cases would require table scans, and the whole point is not to do one. The CASCADE option can be used to automatically
include all dependent tables — but be very careful when using this option, or else you might lose data you did not intend to!

TRUNCATE will not fire any ON DELETE triggers that might exist for the tables. But it will fire ON TRUNCATE triggers. If ON TRUNCATE triggers are defined for any of the tables, then all BEFORE TRUNCATE triggers are fired before any truncation happens, and all AFTER TRUNCATE triggers are fired after the last truncation is performed and any sequences are reset. The triggers will fire in the order that the tables are to be processed (first those listed in the command, and then any that were added due to cascading).

TRUNCATE is not MVCC-safe. After truncation, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the truncation occurred. See Section 13.5 for more details.

TRUNCATE is transaction-safe with respect to the data in the tables: the truncation will be safely rolled back if the surrounding transaction does not commit.

When RESTART IDENTITY is specified, the implied ALTER SEQUENCE RESTART operations are also done transactionally; that is, they will be rolled back if the surrounding transaction does not commit. This is unlike the normal behavior of ALTER SEQUENCE RESTART. Be aware that if any additional sequence operations are done on the restarted sequences before the transaction rolls back, the effects of these operations on the sequences will be rolled back, but not their effects on currval(); that is, after the transaction currval() will continue to reflect the last sequence value obtained inside the failed transaction, even though the sequence itself may no longer be consistent with that. This is similar to the usual behavior of currval() after a failed transaction.

TRUNCATE is not currently supported for foreign tables. This implies that if a specified table has any descendant tables that are foreign, the command will fail.

**Examples**

Truncate the tables bigtable and fattable:

```
TRUNCATE bigtable, fattable;
```

The same, and also reset any associated sequence generators:

```
TRUNCATE bigtable, fattable RESTART IDENTITY;
```

Truncate the table othertable, and cascade to any tables that reference othertable via foreign-key constraints:

```
TRUNCATE othertable CASCADE;
```

**Compatibility**

The SQL:2008 standard includes a TRUNCATE command with the syntax TRUNCATE TABLE tablename. The clauses CONTINUE IDENTITY/RESTART IDENTITY also appear in that standard, but have slightly different though related meanings. Some of the concurrency behavior of this command is left implementation-defined by the standard, so the above notes should be considered and compared with other implementations if necessary.

**See Also**

DELETE
UNLISTEN

UNLISTEN — stop listening for a notification

Synopsis

UNLISTEN { channel | * }

Description

UNLISTEN is used to remove an existing registration for NOTIFY events. UNLISTEN cancels any existing registration of the current PostgreSQL session as a listener on the notification channel named channel. The special wildcard * cancels all listener registrations for the current session.

NOTIFY contains a more extensive discussion of the use of LISTEN and NOTIFY.

Parameters

channel

Name of a notification channel (any identifier).

*

All current listen registrations for this session are cleared.

Notes

You can unlisten something you were not listening for; no warning or error will appear.

At the end of each session, UNLISTEN * is automatically executed.

A transaction that has executed UNLISTEN cannot be prepared for two-phase commit.

Examples

To make a registration:

LISTEN virtual;
NOTIFY virtual;
Asynchronous notification "virtual" received from server process with PID 8448.

Once UNLISTEN has been executed, further NOTIFY messages will be ignored:

UNLISTEN virtual;
NOTIFY virtual;
-- no NOTIFY event is received

Compatibility

There is no UNLISTEN command in the SQL standard.
See Also

LISTEN, NOTIFY
UPDATE

UPDATE — update rows of a table

Synopsis

\[
\begin{align*}
[ \text{WITH} \ [ \text{RECURSIVE} ] \ with\_query \ [, \ ...] ] \\
\text{UPDATE} \ [ \text{ONLY} ] \ table\_name \ [ \ * \ ] \ [ \ [ \ AS \ ] \ alias \ ] \\
\text{SET} \ [ \ \text{column\_name} = \ { \ \text{expression} \ | \ \text{DEFAULT} \ } | \\
( \ \text{column\_name} \ [ , \ ...] ) = \ [ \ \text{ROW} \ ] \ ( \ { \ \text{expression} \ | \\
\text{DEFAULT} \ } [ , \ ...] ) | \\
( \ \text{column\_name} \ [ , \ ...] ) = \ ( \ \text{sub-SELECT} \ ) \\
] \ [ , \ ...] \\
[ \ \text{FROM} \ from\_list \ ] \\
[ \ \text{WHERE} \ \text{condition} \ | \ \text{WHERE CURRENT OF} \ cursor\_name \ ] \\
[ \ \text{RETURING} \ * \ | \ output\_expression \ [ \ [ \ AS \ ] \ output\_name \ ] \\
] \ [ , \ ...] \\
\end{align*}
\]

Description

UPDATE changes the values of the specified columns in all rows that satisfy the condition. Only the columns to be modified need be mentioned in the SET clause; columns not explicitly modified retain their previous values.

There are two ways to modify a table using information contained in other tables in the database: using sub-selects, or specifying additional tables in the FROM clause. Which technique is more appropriate depends on the specific circumstances.

The optional RETURNING clause causes UPDATE to compute and return value(s) based on each row actually updated. Any expression using the table's columns, and/or columns of other tables mentioned in FROM, can be computed. The new (post-update) values of the table's columns are used. The syntax of the RETURNING list is identical to that of the output list of SELECT.

You must have the UPDATE privilege on the table, or at least on the column(s) that are listed to be updated. You must also have the SELECT privilege on any column whose values are read in the expressions or condition.

Parameters

with_query

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the UPDATE query. See Section 7.8 and SELECT for details.

table_name

The name (optionally schema-qualified) of the table to update. If ONLY is specified before the table name, matching rows are updated in the named table only. If ONLY is not specified, matching rows are also updated in any tables inheriting from the named table. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

alias

A substitute name for the target table. When an alias is provided, it completely hides the actual name of the table. For example, given UPDATE foo AS f, the remainder of the UPDATE statement must refer to this table as f not foo.
column_name

The name of a column in the table named by table_name. The column name can be qualified with a subfield name or array subscript, if needed. Do not include the table’s name in the specification of a target column — for example, UPDATE table_name SET table_name.col = 1 is invalid.

expression

An expression to assign to the column. The expression can use the old values of this and other columns in the table.

DEFAULT

Set the column to its default value (which will be NULL if no specific default expression has been assigned to it).

sub-SELECT

A SELECT sub-query that produces as many output columns as are listed in the parenthesized column list preceding it. The sub-query must yield no more than one row when executed. If it yields one row, its column values are assigned to the target columns; if it yields no rows, NULL values are assigned to the target columns. The sub-query can refer to old values of the current row of the table being updated.

from_list

A list of table expressions, allowing columns from other tables to appear in the WHERE condition and the update expressions. This is similar to the list of tables that can be specified in the FROM Clause of a SELECT statement. Note that the target table must not appear in the from_list, unless you intend a self-join (in which case it must appear with an alias in the from_list).

condition

An expression that returns a value of type boolean. Only rows for which this expression returns true will be updated.

cursor_name

The name of the cursor to use in a WHERE CURRENT OF condition. The row to be updated is the one most recently fetched from this cursor. The cursor must be a non-grouping query on the UPDATE’s target table. Note that WHERE CURRENT OF cannot be specified together with a Boolean condition. See DECLARE for more information about using cursors with WHERE CURRENT OF.

output_expression

An expression to be computed and returned by the UPDATE command after each row is updated. The expression can use any column names of the table named by table_name or table(s) listed in FROM. Write * to return all columns.

output_name

A name to use for a returned column.

Outputs

On successful completion, an UPDATE command returns a command tag of the form

UPDATE count
The `count` is the number of rows updated, including matched rows whose values did not change. Note that the number may be less than the number of rows that matched the condition when updates were suppressed by a BEFORE UPDATE trigger. If `count` is 0, no rows were updated by the query (this is not considered an error).

If the UPDATE command contains a RETURNING clause, the result will be similar to that of a SELECT statement containing the columns and values defined in the RETURNING list, computed over the row(s) updated by the command.

**Notes**

When a FROM clause is present, what essentially happens is that the target table is joined to the tables mentioned in the from_list, and each output row of the join represents an update operation for the target table. When using FROM you should ensure that the join produces at most one output row for each row to be modified. In other words, a target row shouldn’t join to more than one row from the other table(s). If it does, then only one of the join rows will be used to update the target row, but which one will be used is not readily predictable.

Because of this indeterminacy, referencing other tables only within sub-selects is safer, though often harder to read and slower than using a join.

In the case of a partitioned table, updating a row might cause it to no longer satisfy the partition constraint of the containing partition. In that case, if there is some other partition in the partition tree for which this row satisfies its partition constraint, then the row is moved to that partition. If there is no such partition, an error will occur. Behind the scenes, the row movement is actually a DELETE and INSERT operation. However, there is a possibility that a concurrent UPDATE or DELETE on the same row may miss this row. For details see the section Section 5.10.2.3. Currently, rows cannot be moved from a partition that is a foreign table to some other partition, but they can be moved into a foreign table if the foreign data wrapper supports it.

**Examples**

Change the word Drama to Dramatic in the column kind of the table films:

```sql
UPDATE films SET kind = 'Dramatic' WHERE kind = 'Drama';
```

Adjust temperature entries and reset precipitation to its default value in one row of the table weather:

```sql
UPDATE weather SET temp_lo = temp_lo+1, temp_hi = temp_lo+15, prcp = DEFAULT
   WHERE city = 'San Francisco' AND date = '2003-07-03';
```

Perform the same operation and return the updated entries:

```sql
UPDATE weather SET temp_lo = temp_lo+1, temp_hi = temp_lo+15, prcp = DEFAULT
   WHERE city = 'San Francisco' AND date = '2003-07-03'
   RETURNING temp_lo, temp_hi, prcp;
```

Use the alternative column-list syntax to do the same update:

```sql
UPDATE weather SET (temp_lo, temp_hi, prcp) = (temp_lo+1, temp_lo +15, DEFAULT)
   WHERE city = 'San Francisco' AND date = '2003-07-03';
```
Increment the sales count of the salesperson who manages the account for Acme Corporation, using the FROM clause syntax:

```
UPDATE employees SET sales_count = sales_count + 1 FROM accounts
    WHERE accounts.name = 'Acme Corporation'
    AND employees.id = accounts.sales_person;
```

Perform the same operation, using a sub-select in the WHERE clause:

```
UPDATE employees SET sales_count = sales_count + 1 WHERE id =
    (SELECT sales_person FROM accounts WHERE name = 'Acme
    Corporation');
```

Update contact names in an accounts table to match the currently assigned salesmen:

```
UPDATE accounts SET (contact_first_name, contact_last_name) =
    (SELECT first_name, last_name FROM salesmen
    WHERE salesmen.id = accounts.sales_id);
```

A similar result could be accomplished with a join:

```
UPDATE accounts SET
    contact_first_name = first_name,
    contact_last_name = last_name
    FROM salesmen WHERE salesmen.id = accounts.sales_id;
```

However, the second query may give unexpected results if `salesmen.id` is not a unique key, whereas the first query is guaranteed to raise an error if there are multiple `id` matches. Also, if there is no match for a particular `accounts.sales_id` entry, the first query will set the corresponding name fields to NULL, whereas the second query will not update that row at all.

Update statistics in a summary table to match the current data:

```
UPDATE summary s SET (sum_x, sum_y, avg_x, avg_y) =
    (SELECT sum(x), sum(y), avg(x), avg(y) FROM data d
    WHERE d.group_id = s.group_id);
```

Attempt to insert a new stock item along with the quantity of stock. If the item already exists, instead update the stock count of the existing item. To do this without failing the entire transaction, use save-points:

```
BEGIN;
    -- other operations
    SAVEPOINT sp1;
    INSERT INTO wines VALUES('Chateau Lafite 2003', '24');
    -- Assume the above fails because of a unique key violation,
    -- so now we issue these commands:
    ROLLBACK TO sp1;
    UPDATE wines SET stock = stock + 24 WHERE winename = 'Chateau
    Lafite 2003';
    -- continue with other operations, and eventually
    COMMIT;
```

Change the `kind` column of the table `films` in the row on which the cursor `c_films` is currently positioned:
UPDATE films SET kind = 'Dramatic' WHERE CURRENT OF c_films;

Compatibility

This command conforms to the SQL standard, except that the FROM and RETURNING clauses are PostgreSQL extensions, as is the ability to use WITH with UPDATE.

Some other database systems offer a FROM option in which the target table is supposed to be listed again within FROM. That is not how PostgreSQL interprets FROM. Be careful when porting applications that use this extension.

According to the standard, the source value for a parenthesized sub-list of target column names can be any row-valued expression yielding the correct number of columns. PostgreSQL only allows the source value to be a row constructor or a sub-SELECT. An individual column's updated value can be specified as DEFAULT in the row-constructor case, but not inside a sub-SELECT.
VACUUM

VACUUM — garbage-collect and optionally analyze a database

Synopsis

VACUUM [ ( option [, ...] ) ] [ table_and_columns [, ...] ]
VACUUM [ FULL ] [ FREEZE ] [ VERBOSE ] [ ANALYZE ]
[ table_and_columns [, ...] ]

where option can be one of:

FULL
FREEZE
VERBOSE
ANALYZE
DISABLE_PAGE_SKIPPING

and table_and_columns is:

table_name [ ( column_name [, ...] ) ]

Description

VACUUM reclaims storage occupied by dead tuples. In normal PostgreSQL operation, tuples that are
deleted or obsoleted by an update are not physically removed from their table; they remain present
until a VACUUM is done. Therefore it's necessary to do VACUUM periodically, especially on frequent-ly-updated tables.

Without a table_and_columns list, VACUUM processes every table and materialized view in the
current database that the current user has permission to vacuum. With a list, VACUUM processes only
those table(s).

VACUUM ANALYZE performs a VACUUM and then an ANALYZE for each selected table. This is a
handy combination form for routine maintenance scripts. See ANALYZE for more details about its
processing.

Plain VACUUM (without FULL) simply reclaims space and makes it available for re-use. This form of
the command can operate in parallel with normal reading and writing of the table, as an exclusive lock
is not obtained. However, extra space is not returned to the operating system (in most cases); it's just
kept available for re-use within the same table. VACUUM FULL rewrites the entire contents of the table
into a new disk file with no extra space, allowing unused space to be returned to the operating system.
This form is much slower and requires an exclusive lock on each table while it is being processed.

When the option list is surrounded by parentheses, the options can be written in any order. Without
parentheses, options must be specified in exactly the order shown above. The parenthesized syntax
was added in PostgreSQL 9.0; the unparenthesized syntax is deprecated.

Parameters

FULL

Selects “full” vacuum, which can reclaim more space, but takes much longer and exclusively
locks the table. This method also requires extra disk space, since it writes a new copy of the table.
and doesn't release the old copy until the operation is complete. Usually this should only be used when a significant amount of space needs to be reclaimed from within the table.

**FREEZE**

Selects aggressive “freezing” of tuples. Specifying **FREEZE** is equivalent to performing **VACUUM** with the vacuum_freeze_min_age and vacuum_freeze_table_age parameters set to zero. Aggressive freezing is always performed when the table is rewritten, so this option is redundant when **FULL** is specified.

**VERBOSE**

Prints a detailed vacuum activity report for each table.

**ANALYZE**

Updates statistics used by the planner to determine the most efficient way to execute a query.

**DISABLE_PAGE_SKIPPING**

Normally, **VACUUM** will skip pages based on the visibility map. Pages where all tuples are known to be frozen can always be skipped, and those where all tuples are known to be visible to all transactions may be skipped except when performing an aggressive vacuum. Furthermore, except when performing an aggressive vacuum, some pages may be skipped in order to avoid waiting for other sessions to finish using them. This option disables all page-skipping behavior, and is intended to be used only when the contents of the visibility map are suspect, which should happen only if there is a hardware or software issue causing database corruption.

**table_name**

The name (optionally schema-qualified) of a specific table or materialized view to vacuum. If the specified table is a partitioned table, all of its leaf partitions are vacuumed.

**column_name**

The name of a specific column to analyze. Defaults to all columns. If a column list is specified, **ANALYZE** must also be specified.

**Outputs**

When **VERBOSE** is specified, **VACUUM** emits progress messages to indicate which table is currently being processed. Various statistics about the tables are printed as well.

**Notes**

To vacuum a table, one must ordinarily be the table's owner or a superuser. However, database owners are allowed to vacuum all tables in their databases, except shared catalogs. (The restriction for shared catalogs means that a true database-wide **VACUUM** can only be performed by a superuser.) **VACUUM** will skip over any tables that the calling user does not have permission to vacuum.

**VACUUM** cannot be executed inside a transaction block.

For tables with GIN indexes, **VACUUM** (in any form) also completes any pending index insertions, by moving pending index entries to the appropriate places in the main GIN index structure. See Section 66.4.1 for details.

We recommend that active production databases be vacuumed frequently (at least nightly), in order to remove dead rows. After adding or deleting a large number of rows, it might be a good idea to issue a **VACUUM ANALYZE** command for the affected table. This will update the system catalogs.
with the results of all recent changes, and allow the PostgreSQL query planner to make better choices in planning queries.

The **FULL** option is not recommended for routine use, but might be useful in special cases. An example is when you have deleted or updated most of the rows in a table and would like the table to physically shrink to occupy less disk space and allow faster table scans. **VACUUM** **FULL** will usually shrink the table more than a plain **VACUUM** would.

**VACUUM** causes a substantial increase in I/O traffic, which might cause poor performance for other active sessions. Therefore, it is sometimes advisable to use the cost-based vacuum delay feature. See Section 19.4.4 for details.

PostgreSQL includes an “autovacuum” facility which can automate routine vacuum maintenance. For more information about automatic and manual vacuuming, see Section 24.1.

**Examples**

To clean a single table `onek`, analyze it for the optimizer and print a detailed vacuum activity report:

```
VACUUM (VERBOSE, ANALYZE) onek;
```

**Compatibility**

There is no **VACUUM** statement in the SQL standard.

**See Also**

vacuumdb, Section 19.4.4, Section 24.1.6
VALUES

VALUES — compute a set of rows

Synopsis

VALUES ( expression [, ...] ) [, ...]
[ ORDER BY sort_expression [ ASC | DESC | USING operator ]
[ , ... ] ]
[ LIMIT { count | ALL } ]
[ OFFSET start [ ROW | ROWS ] ]
[ FETCH { FIRST | NEXT } [ count ] { ROW | ROWS } ONLY ]

Description

VALUES computes a row value or set of row values specified by value expressions. It is most commonly used to generate a “constant table” within a larger command, but it can be used on its own.

When more than one row is specified, all the rows must have the same number of elements. The data types of the resulting table’s columns are determined by combining the explicit or inferred types of the expressions appearing in that column, using the same rules as for UNION (see Section 10.5).

Within larger commands, VALUES is syntactically allowed anywhere that SELECT is. Because it is treated like a SELECT by the grammar, it is possible to use the ORDER BY, LIMIT (or equivalently FETCH FIRST), and OFFSET clauses with a VALUES command.

Parameters

expression

A constant or expression to compute and insert at the indicated place in the resulting table (set of rows). In a VALUES list appearing at the top level of an INSERT, an expression can be replaced by DEFAULT to indicate that the destination column’s default value should be inserted. DEFAULT cannot be used when VALUES appears in other contexts.

sort_expression

An expression or integer constant indicating how to sort the result rows. This expression can refer to the columns of the VALUES result as column1, column2, etc. For more details see ORDER BY Clause.

operator

A sorting operator. For details see ORDER BY Clause.

count

The maximum number of rows to return. For details see LIMIT Clause.

start

The number of rows to skip before starting to return rows. For details see LIMIT Clause.

Notes

VALUES lists with very large numbers of rows should be avoided, as you might encounter out-of-memory failures or poor performance. VALUES appearing within INSERT is a special case (because
the desired column types are known from the INSERT's target table, and need not be inferred by scanning the VALUES list), so it can handle larger lists than are practical in other contexts.

Examples

A bare VALUES command:

VALUES (1, 'one'), (2, 'two'), (3, 'three');

This will return a table of two columns and three rows. It's effectively equivalent to:

SELECT 1 AS column1, 'one' AS column2
UNION ALL
SELECT 2, 'two'
UNION ALL
SELECT 3, 'three';

More usually, VALUES is used within a larger SQL command. The most common use is in INSERT:

INSERT INTO films (code, title, did, date_prod, kind)
VALUES ('T_601', 'Yojimbo', 106, '1961-06-16', 'Drama');

In the context of INSERT, entries of a VALUES list can be DEFAULT to indicate that the column default should be used here instead of specifying a value:

INSERT INTO films VALUES
 ('UA502', 'Bananas', 105, DEFAULT, 'Comedy', '82 minutes'),
 ('T_601', 'Yojimbo', 106, DEFAULT, 'Drama', DEFAULT);

VALUES can also be used where a sub-SELECT might be written, for example in a FROM clause:

SELECT f.*
 FROM films f, (VALUES('MGM', 'Horror'), ('UA', 'Sci-Fi')) AS t (studio, kind)
 WHERE f.studio = t.studio AND f.kind = t.kind;

UPDATE employees SET salary = salary * v.increase
 FROM (VALUES(1, 200000, 1.2), (2, 400000, 1.4)) AS v (depno, target, increase)
 WHERE employees.depno = v.depno AND employees.sales >= v.target;

Note that an AS clause is required when VALUES is used in a FROM clause, just as is true for SELECT. It is not required that the AS clause specify names for all the columns, but it's good practice to do so. (The default column names for VALUES are column1, column2, etc in PostgreSQL, but these names might be different in other database systems.)

When VALUES is used in INSERT, the values are all automatically coerced to the data type of the corresponding destination column. When it's used in other contexts, it might be necessary to specify the correct data type. If the entries are all quoted literal constants, coercing the first is sufficient to determine the assumed type for all:

SELECT * FROM machines
WHERE ip_address IN (VALUES('192.168.0.1'::inet), ('192.168.0.10'), ('192.168.1.43'));
VALUES

Tip

For simple IN tests, it's better to rely on the list-of-scalars form of IN than to write a VALUES query as shown above. The list of scalars method requires less writing and is often more efficient.

Compatibility

VALUES conforms to the SQL standard. LIMIT and OFFSET are PostgreSQL extensions; see also under SELECT.

See Also

INSERT, SELECT
PostgreSQL Client Applications

This part contains reference information for PostgreSQL client applications and utilities. Not all of these commands are of general utility; some might require special privileges. The common feature of these applications is that they can be run on any host, independent of where the database server resides.

When specified on the command line, user and database names have their case preserved — the presence of spaces or special characters might require quoting. Table names and other identifiers do not have their case preserved, except where documented, and might require quoting.

Table of Contents

clusterdb ................................................................. 1817
createdb ................................................................. 1820
createuser ............................................................... 1823
dropdb ................................................................. 1827
dropuser ............................................................... 1830
ecpg .......................................................................... 1833
pg_basebackup ............................................................. 1835
pgbench ................................................................. 1842
pg_config ................................................................. 1858
pg_dump ................................................................. 1861
pg_dumpall ................................................................. 1873
pg_isready ................................................................. 1879
pg_receivewal .............................................................. 1881
pg_recvlogical ............................................................ 1885
pg_restore .............................................................. 1889
psql ........................................................................ 1898
reindexdb ................................................................. 1937
vacuumdb ................................................................. 1940
clusterdb

clusterdb — cluster a PostgreSQL database

Synopsis

clusterdb [connection-option...] [--verbose | -v] [--table | -t table] ... [dbname]
clusterdb [connection-option...] [--verbose | -v] --all | -a

Description

clusterdb is a utility for reclustering tables in a PostgreSQL database. It finds tables that have previously been clustered, and clusters them again on the same index that was last used. Tables that have never been clustered are not affected.

clusterdb is a wrapper around the SQL command CLUSTER. There is no effective difference between clustering databases via this utility and via other methods for accessing the server.

Options

clusterdb accepts the following command-line arguments:

-a
--all

Cluster all databases.

[-d] dbname
[--dbname=] dbname

Specifies the name of the database to be clustered. If this is not specified and -a (or --all) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

-e
--echo

Echo the commands that clusterdb generates and sends to the server.

-q
--quiet

Do not display progress messages.

-t table
--table=table

Cluster table only. Multiple tables can be clustered by writing multiple -t switches.

-v
--verbose

Print detailed information during processing.
-V
--version

Print the clusterdb version and exit.

-?
--help

Show help about clusterdb command line arguments, and exit.

clusterdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force clusterdb to prompt for a password before connecting to a database.

This option is never essential, since clusterdb will automatically prompt for a password if the server demands password authentication. However, clusterdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to discover what other databases should be clustered. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.

Environment

PGDATABASE
PGHOST
PGPORT
PGUSER

Default connection parameters
This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

**Diagnostics**

In case of difficulty, see CLUSTER and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

**Examples**

To cluster the database *test*:

```
$ clusterdb test
```

To cluster a single table *foo* in a database named *xyzzy*:

```
$ clusterdb --table=foo xyzzy
```

**See Also**

CLUSTER
createdb

createdb — create a new PostgreSQL database

Synopsis

createdb [connection-option...] [option...] [dbname [description]]

Description

createdb creates a new PostgreSQL database.

Normally, the database user who executes this command becomes the owner of the new database. However, a different owner can be specified via the -O option, if the executing user has appropriate privileges.

createdb is a wrapper around the SQL command CREATE DATABASE. There is no effective difference between creating databases via this utility and via other methods for accessing the server.

Options

createdb accepts the following command-line arguments:

dbname

  Specifies the name of the database to be created. The name must be unique among all PostgreSQL databases in this cluster. The default is to create a database with the same name as the current system user.

description

  Specifies a comment to be associated with the newly created database.

-D tablespace
  --tablespace=tablespace

  Specifies the default tablespace for the database. (This name is processed as a double-quoted identifier.)

-e
  --echo

  Echo the commands that createdb generates and sends to the server.

-E encoding
  --encoding=encoding

  Specifies the character encoding scheme to be used in this database. The character sets supported by the PostgreSQL server are described in Section 23.3.1.

-l locale
  --locale=locale

  Specifies the locale to be used in this database. This is equivalent to specifying both --lc-collate and --lc-ctype.

  --lc-collate=locale

  Specifies the LC_COLLATE setting to be used in this database.
createdb

--lc-ctype=locale

Specifies the LC_CTYPE setting to be used in this database.

-O owner
--owner=owner

Specifies the database user who will own the new database. (This name is processed as a double-quoted identifier.)

-T template
--template=template

Specifies the template database from which to build this database. (This name is processed as a double-quoted identifier.)

-V
--version

Print the createdb version and exit.

-?
--help

Show help about createdb command line arguments, and exit.

The options --D, --l, --E, --O, and --T correspond to options of the underlying SQL command CREATE DATABASE; see there for more information about them.

createdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or the local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force createdb to prompt for a password before connecting to a database.

This option is never essential, since createdb will automatically prompt for a password if the server demands password authentication. However, createdb will waste a connection attempt finding out
that the server wants a password. In some cases it is worth typing `--no-password` to avoid the extra connection attempt.

`--maintenance-db=dbname`

Specifies the name of the database to connect to when creating the new database. If not specified, the `postgres` database will be used; if that does not exist (or if it is the name of the new database being created), `template1` will be used.

**Environment**

`PGDATABASE`

If set, the name of the database to create, unless overridden on the command line.

`PGHOST`  
`PGPORT`  
`PGUSER`

Default connection parameters. `PGUSER` also determines the name of the database to create, if it is not specified on the command line or by `PGDATABASE`.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by `libpq` (see Section 34.14).

**Diagnostics**

In case of difficulty, see `CREATE DATABASE` and `psql` for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the `libpq` front-end library will apply.

**Examples**

To create the database `demo` using the default database server:

```
$ createdb demo
```

To create the database `demo` using the server on host `eden`, port 5000, using the `template0` template database, here is the command-line command and the underlying SQL command:

```
$ createdb --host=eden --port=5000 --template=template0 --dbname demo
```

```
CREATE DATABASE demo TEMPLATE template0;
```

**See Also**

`dropdb`, `CREATE DATABASE`
createuser

createuser — define a new PostgreSQL user account

Synopsis

createuser [connection-option...][option...][username]

Description

createuser creates a new PostgreSQL user (or more precisely, a role). Only superusers and users with
CREATEROLE privilege can create new users, so createuser must be invoked by someone who can
connect as a superuser or a user with CREATEROLE privilege.

If you wish to create a new superuser, you must connect as a superuser, not merely with CREATEROLE
privilege. Being a superuser implies the ability to bypass all access permission checks within the
database, so superuserdom should not be granted lightly.

createuser is a wrapper around the SQL command CREATE ROLE. There is no effective difference
between creating users via this utility and via other methods for accessing the server.

Options

createuser accepts the following command-line arguments:

username

Specifies the name of the PostgreSQL user to be created. This name must be different from all
existing roles in this PostgreSQL installation.

-c number
--connection-limit=number

Set a maximum number of connections for the new user. The default is to set no limit.

-d
--createdb

The new user will be allowed to create databases.

-D
--no-createdb

The new user will not be allowed to create databases. This is the default.

-e
--echo

Echo the commands that createuser generates and sends to the server.

-E
--encrypted

This option is obsolete but still accepted for backward compatibility.

-g role
--role=role

Indicates role to which this role will be added immediately as a new member. Multiple roles to
which this role will be added as a member can be specified by writing multiple -g switches.
createuser

-i
--inherit

The new role will automatically inherit privileges of roles it is a member of. This is the default.

-I
--no-inherit

The new role will not automatically inherit privileges of roles it is a member of.

--interactive

Prompt for the user name if none is specified on the command line, and also prompt for whichever of the options -d/-D, -r/-R, -s/-S is not specified on the command line. (This was the default behavior up to PostgreSQL 9.1.)

-l
--login

The new user will be allowed to log in (that is, the user name can be used as the initial session user identifier). This is the default.

-L
--no-login

The new user will not be allowed to log in. (A role without login privilege is still useful as a means of managing database permissions.)

-P
--pwprompt

If given, createuser will issue a prompt for the password of the new user. This is not necessary if you do not plan on using password authentication.

-r
--createrole

The new user will be allowed to create new roles (that is, this user will have CREATEROLE privilege).

-R
--no-createrole

The new user will not be allowed to create new roles. This is the default.

-s
--superuser

The new user will be a superuser.

-S
--no-superuser

The new user will not be a superuser. This is the default.

-V
--version

Print the createuser version and exit.
The new user will have the REPLICATION privilege, which is described more fully in the documentation for CREATE ROLE.

The new user will not have the REPLICATION privilege, which is described more fully in the documentation for CREATE ROLE.

Show help about createuser command line arguments, and exit.

createuser also accepts the following command-line arguments for connection parameters:

```
-h host
--host=host
```

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

```
-p port
--port=port
```

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

```
-U username
--username=username
```

User name to connect as (not the user name to create).

```
-w
--no-password
```

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

```
-W
--password
```

Force createuser to prompt for a password (for connecting to the server, not for the password of the new user).

This option is never essential, since createuser will automatically prompt for a password if the server demands password authentication. However, createuser will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

**Environment**

PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).
Diagnostics

In case of difficulty, see CREATE ROLE and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Examples

To create a user `joe` on the default database server:

```
$ createuser joe
```

To create a user `joe` on the default database server with prompting for some additional attributes:

```
$ createuser --interactive joe
```

Shall the new role be a superuser? (y/n) n
Shall the new role be allowed to create databases? (y/n) n
Shall the new role be allowed to create more new roles? (y/n) n

To create the same user `joe` using the server on host `eden`, port 5000, with attributes explicitly specified, taking a look at the underlying command:

```
$ createuser -h eden -p 5000 -S -D -R -e joe
CREATE ROLE joe NOSUPERUSER NOCREATEDB NOCREATEROLE INHERIT LOGIN;
```

To create the user `joe` as a superuser, and assign a password immediately:

```
$ createuser -P -s -e joe
```

Enter password for new role: `xyzzy`
Enter it again: `xyzzy`

```
CREATE ROLE joe PASSWORD 'md5b5f5ba1a423792b526f799ae4eb3d59e'
  SUPERUSER CREATEDB CREATEROLE INHERIT LOGIN;
```

In the above example, the new password isn’t actually echoed when typed, but we show what was typed for clarity. As you see, the password is encrypted before it is sent to the client.

See Also

`dropuser`, `CREATE ROLE`
dropdb

dropdb — remove a PostgreSQL database

Synopsis

dropdb [connection-option...][option... ] dbname

Description

dropdb destroys an existing PostgreSQL database. The user who executes this command must be a
database superuser or the owner of the database.

dropdb is a wrapper around the SQL command DROP DATABASE. There is no effective difference
between dropping databases via this utility and via other methods for accessing the server.

Options

dropdb accepts the following command-line arguments:

dbname

Specifies the name of the database to be removed.

-e
--echo

Echo the commands that dropdb generates and sends to the server.

-i
--interactive

Issues a verification prompt before doing anything destructive.

-V
--version

Print the dropdb version and exit.

--if-exists

Do not throw an error if the database does not exist. A notice is issued in this case.

-?
--help

Show help about dropdb command line arguments, and exit.

dropdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with
a slash, it is used as the directory for the Unix domain socket.
-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force dropdb to prompt for a password before connecting to a database.

This option is never essential, since dropdb will automatically prompt for a password if the server demands password authentication. However, dropdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to in order to drop the target database. If not specified, the postgres database will be used; if that does not exist (or is the database being dropped), template1 will be used.

Environment

PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Diagnostics

In case of difficulty, see DROP DATABASE and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Examples

To destroy the database demo on the default database server:

$ dropdb demo

To destroy the database demo using the server on host eden, port 5000, with verification and a peek at the underlying command:


See Also

createdb, DROP DATABASE
dropuser

dropuser — remove a PostgreSQL user account

Synopsis

dropuser [connection-option...] [option...] [username]

Description

dropuser removes an existing PostgreSQL user. Only superusers and users with the CREATEROLE privilege can remove PostgreSQL users. (To remove a superuser, you must yourself be a superuser.)

dropuser is a wrapper around the SQL command DROP ROLE. There is no effective difference between dropping users via this utility and via other methods for accessing the server.

Options

dropuser accepts the following command-line arguments:

username

Specifies the name of the PostgreSQL user to be removed. You will be prompted for a name if none is specified on the command line and the -i/--interactive option is used.

-e
--echo

Echo the commands that dropuser generates and sends to the server.

-i
--interactive

Prompt for confirmation before actually removing the user, and prompt for the user name if none is specified on the command line.

-V
--version

Print the dropuser version and exit.

--if-exists

Do not throw an error if the user does not exist. A notice is issued in this case.

-?
--help

Show help about dropuser command line arguments, and exit.

dropuser also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.
-p port  
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username  
--username=username

User name to connect as (not the user name to drop).

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force dropuser to prompt for a password before connecting to a database.

This option is never essential, since dropuser will automatically prompt for a password if the server demands password authentication. However, dropuser will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

Environment

PGHOST  
PGPORT  
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Diagnostics

In case of difficulty, see DROP ROLE and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Examples

To remove user joe from the default database server:

$ dropuser joe

To remove user joe using the server on host eden, port 5000, with verification and a peek at the underlying command:

$ dropuser -p 5000 -h eden -i -e joe
Role "joe" will be permanently removed.  
Are you sure? (y/n) y
DROP ROLE joe;

See Also
createuser, DROP ROLE
ecpg

ecpg — embedded SQL C preprocessor

Synopsis

ecpg [option...] file...

Description

ecpg is the embedded SQL preprocessor for C programs. It converts C programs with embedded SQL statements to normal C code by replacing the SQL invocations with special function calls. The output files can then be processed with any C compiler tool chain.

ecpg will convert each input file given on the command line to the corresponding C output file. Input files preferably have the extension .pgc. The extension will be replaced by .c to determine the output file name. The output file name can also be overridden using the -o option.

This reference page does not describe the embedded SQL language. See Chapter 36 for more information on that topic.

Options

ecpg accepts the following command-line arguments:

- c

  Automatically generate certain C code from SQL code. Currently, this works for EXEC SQL TYPE.

- C mode

  Set a compatibility mode. mode can be INFORMIX, INFORMIX_SE, or ORACLE.

- D symbol

  Define a C preprocessor symbol.

- i

  Parse system include files as well.

- I directory

  Specify an additional include path, used to find files included via EXEC SQL INCLUDE. Defaults are . (current directory), /usr/local/include, the PostgreSQL include directory which is defined at compile time (default: /usr/local/pgsql/include), and /usr/include, in that order.

- o filename

  Specifies that ecpg should write all its output to the given filename.

- r option

  Selects run-time behavior. Option can be one of the following:
no_indicator

Do not use indicators but instead use special values to represent null values. Historically there have been databases using this approach.

prepare

Prepare all statements before using them. Libecpg will keep a cache of prepared statements and reuse a statement if it gets executed again. If the cache runs full, libecpg will free the least used statement.

questionmarks

Allow question mark as placeholder for compatibility reasons. This used to be the default long ago.

-t

Turn on autocommit of transactions. In this mode, each SQL command is automatically committed unless it is inside an explicit transaction block. In the default mode, commands are committed only when EXEC SQL COMMIT is issued.

-v

Print additional information including the version and the "include" path.

--version

Print the ecpg version and exit.

-?

--help

Show help about ecpg command line arguments, and exit.

Notes

When compiling the preprocessed C code files, the compiler needs to be able to find the ECPG header files in the PostgreSQL include directory. Therefore, you might have to use the -I option when invoking the compiler (e.g., -I/usr/local/pgsql/include).

Programs using C code with embedded SQL have to be linked against the libecpg library, for example using the linker options -L/usr/local/pgsql/lib -lecpg.

The value of either of these directories that is appropriate for the installation can be found out using pg_config.

Examples

If you have an embedded SQL C source file named prog1.pgc, you can create an executable program using the following sequence of commands:

ecpg prog1.pgc
cc -I/usr/local/pgsql/include -c prog1.c
cc -o prog1 prog1.o -L/usr/local/pgsql/lib -lecpg
**Description**

pg_basebackup is used to take base backups of a running PostgreSQL database cluster. These backups are taken without affecting other clients to the database, and can be used both for point-in-time recovery (see Section 25.3) and as the starting point for a log shipping or streaming replication standby servers (see Section 26.2).

pg_basebackup makes a binary copy of the database cluster files, while making sure the system is put in and out of backup mode automatically. Backups are always taken of the entire database cluster; it is not possible to back up individual databases or database objects. For individual database backups, a tool such as pg_dump must be used.

The backup is made over a regular PostgreSQL connection, and uses the replication protocol. The connection must be made with a superuser or a user having `REPLICATION` permissions (see Section 21.2), and `pg_hba.conf` must explicitly permit the replication connection. The server must also be configured with `max_wal_senders` set high enough to leave at least one session available for the backup and one for WAL streaming (if used).

There can be multiple pg_basebackups running at the same time, but it is better from a performance point of view to take only one backup, and copy the result.

pg_basebackup can make a base backup from not only the master but also the standby. To take a backup from the standby, set up the standby so that it can accept replication connections (that is, set `max_wal_senders` and `hot_standby`, and configure host-based authentication). You will also need to enable `full_page_writes` on the master.

Note that there are some limitations in an online backup from the standby:

- The backup history file is not created in the database cluster backed up.
- If you are using `-X none`, there is no guarantee that all WAL files required for the backup are archived at the end of backup.
- If the standby is promoted to the master during online backup, the backup fails.
- All WAL records required for the backup must contain sufficient full-page writes, which requires you to enable `full_page_writes` on the master and not to use a tool like `pg_compresslog` as `archive_command` to remove full-page writes from WAL files.

**Options**

The following command-line options control the location and format of the output.

```
-D directory
--pgdata=directory
```

Directory to write the output to. pg_basebackup will create the directory and any parent directories if necessary. The directory may already exist, but it is an error if the directory already exists and is not empty.
When the backup is in tar mode, and the directory is specified as - (dash), the tar file will be written to stdout.

This option is required.

- \texttt{-F format} \hfill  --format=format

Selects the format for the output. \texttt{format} can be one of the following:

\begin{itemize}
  \item \texttt{p} \hfill \texttt{plain}
    \begin{itemize}
    \item Write the output as plain files, with the same layout as the current data directory and tablespaces. When the cluster has no additional tablespaces, the whole database will be placed in the target directory. If the cluster contains additional tablespaces, the main data directory will be placed in the target directory, but all other tablespaces will be placed in the same absolute path as they have on the server.
    \end{itemize}
  \item \texttt{t} \hfill \texttt{tar}
    \begin{itemize}
    \item Write the output as tar files in the target directory. The main data directory will be written to a file named \texttt{base.tar}, and all other tablespaces will be named after the tablespace OID.
    \end{itemize}
\end{itemize}

If the value - (dash) is specified as target directory, the tar contents will be written to standard output, suitable for piping to for example gzip. This is only possible if the cluster has no additional tablespaces and WAL streaming is not used.

- \texttt{-r rate} \hfill  --max-rate=rate

The maximum transfer rate of data transferred from the server. Values are in kilobytes per second. Use a suffix of \texttt{M} to indicate megabytes per second. A suffix of \texttt{k} is also accepted, and has no effect. Valid values are between 32 kilobytes per second and 1024 megabytes per second.

The purpose is to limit the impact of \texttt{pg_basebackup} on the running server.

This option always affects transfer of the data directory. Transfer of WAL files is only affected if the collection method is \texttt{fetch}.

- \texttt{-R} \hfill  --write-recovery-conf

Write a minimal \texttt{recovery.conf} in the output directory (or into the base archive file when using tar format) to ease setting up a standby server. The \texttt{recovery.conf} file will record the connection settings and, if specified, the replication slot that \texttt{pg_basebackup} is using, so that the streaming replication will use the same settings later on.

- \texttt{-T olddir=newdir} \hfill  --tablespace-mapping=olddir=newdir

Relocate the tablespace in directory \texttt{olddir} to \texttt{newdir} during the backup. To be effective, \texttt{olddir} must exactly match the path specification of the tablespace as it is currently defined. (But it is not an error if there is no tablespace in \texttt{olddir} contained in the backup.) Both \texttt{olddir} and \texttt{newdir} must be absolute paths. If a path happens to contain a \texttt{=} sign, escape it with a backslash. This option can be specified multiple times for multiple tablespaces. See examples below.

If a tablespace is relocated in this way, the symbolic links inside the main data directory are updated to point to the new location. So the new data directory is ready to be used for a new server instance with all tablespaces in the updated locations.
--waldir=waldir

Specifies the location for the write-ahead log directory. waldir must be an absolute path. The write-ahead log directory can only be specified when the backup is in plain mode.

-X method
--wal-method=method

Includes the required write-ahead log files (WAL files) in the backup. This will include all write-ahead logs generated during the backup. Unless the method none is specified, it is possible to start a postmaster directly in the extracted directory without the need to consult the log archive, thus making this a completely standalone backup.

The following methods for collecting the write-ahead logs are supported:

n
none

Don't include write-ahead log in the backup.

f
fetch

The write-ahead log files are collected at the end of the backup. Therefore, it is necessary for the wal_keep_segments parameter to be set high enough that the log is not removed before the end of the backup. If the log has been rotated when it's time to transfer it, the backup will fail and be unusable.

The write-ahead log files will be written to the base.tar file.

s
stream

Stream the write-ahead log while the backup is created. This will open a second connection to the server and start streaming the write-ahead log in parallel while running the backup. Therefore, it will use up two connections configured by the max_wal_senders parameter. As long as the client can keep up with write-ahead log received, using this mode requires no extra write-ahead logs to be saved on the master.

The write-ahead log files are written to a separate file named pg_wal.tar (if the server is a version earlier than 10, the file will be named pg_xlog.tar).

This value is the default.

-z
--gzip

Enables gzip compression of tar file output, with the default compression level. Compression is only available when using the tar format, and the suffix .gz will automatically be added to all tar filenames.

-Z level
--compress=level

Enables gzip compression of tar file output, and specifies the compression level (0 through 9, 0 being no compression and 9 being best compression). Compression is only available when using the tar format, and the suffix .gz will automatically be added to all tar filenames.

The following command-line options control the generation of the backup and the running of the program.
-c fast|spread
--checkpoint=fast|spread

Sets checkpoint mode to fast (immediate) or spread (default) (see Section 25.3.3).

-C
--create-slot

This option causes creation of a replication slot named by the --slot option before starting the backup. An error is raised if the slot already exists.

-l label
--label=label

Sets the label for the backup. If none is specified, a default value of "pg_basebackup base backup" will be used.

-n
--no-clean

By default, when pg_basebackup aborts with an error, it removes any directories it might have created before discovering that it cannot finish the job (for example, data directory and write-ahead log directory). This option inhibits tidying-up and is thus useful for debugging.

Note that tablespace directories are not cleaned up either way.

-N
--no-sync

By default, pg_basebackup will wait for all files to be written safely to disk. This option causes pg_basebackup to return without waiting, which is faster, but means that a subsequent operating system crash can leave the base backup corrupt. Generally, this option is useful for testing but should not be used when creating a production installation.

-P
--progress

Enables progress reporting. Turning this on will deliver an approximate progress report during the backup. Since the database may change during the backup, this is only an approximation and may not end at exactly 100%. In particular, when WAL log is included in the backup, the total amount of data cannot be estimated in advance, and in this case the estimated target size will increase once it passes the total estimate without WAL.

When this is enabled, the backup will start by enumerating the size of the entire database, and then go back and send the actual contents. This may make the backup take slightly longer, and in particular it will take longer before the first data is sent.

-S slotname
--slot=slotname

This option can only be used together with -X stream. It causes the WAL streaming to use the specified replication slot. If the base backup is intended to be used as a streaming replication standby using replication slots, it should then use the same replication slot name in recovery.conf. That way, it is ensured that the server does not remove any necessary WAL data in the time between the end of the base backup and the start of streaming replication.

The specified replication slot has to exist unless the option -C is also used.

If this option is not specified and the server supports temporary replication slots (version 10 and later), then a temporary replication slot is automatically used for WAL streaming.
-v
--verbose

Enables verbose mode. Will output some extra steps during startup and shutdown, as well as show
the exact file name that is currently being processed if progress reporting is also enabled.

--no-slot

This option prevents the creation of a temporary replication slot during the backup even if it's
supported by the server.

Temporary replication slots are created by default if no slot name is given with the option -S
when using log streaming.

The main purpose of this option is to allow taking a base backup when the server is out of free
replication slots. Using replication slots is almost always preferred, because it prevents needed
WAL from being removed by the server during the backup.

--no-verify-checksums

Disables verification of checksums, if they are enabled on the server the base backup is taken from.

By default, checksums are verified and checksum failures will result in a non-zero exit status.
However, the base backup will not be removed in such a case, as if the --no-clean option
had been used.

The following command-line options control the database connection parameters.

-d connstr
--dbname=connstr

Specifies parameters used to connect to the server, as a connection string. See Section 34.1.1  for
more information.

The option is called --dbname for consistency with other client applications, but because
pg_basebackup doesn't connect to any particular database in the cluster, database name in the
connection string will be ignored.

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a
slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST
environment variable, if set, else a Unix domain socket connection is attempted.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening
for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-s interval
--status-interval=interval

Specifies the number of seconds between status packets sent back to the server. This allows for
easier monitoring of the progress from server. A value of zero disables the periodic status updates
completely, although an update will still be sent when requested by the server, to avoid timeout
disconnect. The default value is 10 seconds.

-U username
--username=username

User name to connect as.
-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_basebackup to prompt for a password before connecting to a database.

This option is never essential, since pg_basebackup will automatically prompt for a password if
the server demands password authentication. However, pg_basebackup will waste a connection
attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid
the extra connection attempt.

Other options are also available:

-V
--version

Print the pg_basebackup version and exit.

-?
--help

Show help about pg_basebackup command line arguments, and exit.

Environment

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq
(see Section 34.14).

Notes

At the beginning of the backup, a checkpoint needs to be written on the server the backup is taken
from. Especially if the option --checkpoint=fast is not used, this can take some time during
which pg_basebackup will appear to be idle.

The backup will include all files in the data directory and tablespaces, including the configuration files
and any additional files placed in the directory by third parties, except certain temporary files managed
by PostgreSQL. But only regular files and directories are copied, except that symbolic links used for
tablespaces are preserved. Symbolic links pointing to certain directories known to PostgreSQL are
copied as empty directories. Other symbolic links and special device files are skipped. See Section 53.4
for the precise details.

Tablespaces will in plain format by default be backed up to the same path they have on the server,
unless the option --tablespace-mapping is used. Without this option, running a plain format
base backup on the same host as the server will not work if tablespaces are in use, because the backup
would have to be written to the same directory locations as the original tablespaces.

When tar format mode is used, it is the user’s responsibility to unpack each tar file before starting the
PostgreSQL server. If there are additional tablespaces, the tar files for them need to be unpacked in
the correct locations. In this case the symbolic links for those tablespaces will be created by the server
according to the contents of the tablespace_map file that is included in the base.tar file.

pg_basebackup works with servers of the same or an older major version, down to 9.1. However,
WAL streaming mode (-X stream) only works with server version 9.3 and later, and tar format
mode (--format=tar) of the current version only works with server version 9.5 or later.
pg_basebackup will preserve group permissions in both the plain and tar formats if group permissions are enabled on the source cluster.

Examples

To create a base backup of the server at mydbserver and store it in the local directory /usr/local/pgsql/data:

$ pg_basebackup -h mydbserver -D /usr/local/pgsql/data

To create a backup of the local server with one compressed tar file for each tablespace, and store it in the directory backup, showing a progress report while running:

$ pg_basebackup -D backup -Ft -z -P

To create a backup of a single-tablespace local database and compress this with bzip2:

$ pg_basebackup -D - -Ft -X fetch | bzip2 > backup.tar.bz2

(This command will fail if there are multiple tablespaces in the database.)

To create a backup of a local database where the tablespace in /opt/ts is relocated to .//back-up/ts:

$ pg_basebackup -D backup/data -T /opt/ts=$(pwd)/backup/ts

See Also

pg_dump
pgbench

pgbench — run a benchmark test on PostgreSQL

Synopsis

pgbench -i [option...] [dbname]
pgbench [option...] [dbname]

Description

pgbench is a simple program for running benchmark tests on PostgreSQL. It runs the same sequence of SQL commands over and over, possibly in multiple concurrent database sessions, and then calculates the average transaction rate (transactions per second). By default, pgbench tests a scenario that is loosely based on TPC-B, involving five SELECT, UPDATE, and INSERT commands per transaction. However, it is easy to test other cases by writing your own transaction script files.

Typical output from pgbench looks like:

transaction type: <builtin: TPC-B (sort of)>
scaling factor: 10
query mode: simple
number of clients: 10
number of threads: 1
number of transactions per client: 1000
number of transactions actually processed: 10000/10000
tps = 85.184871 (including connections establishing)
tps = 85.296346 (excluding connections establishing)

The first six lines report some of the most important parameter settings. The next line reports the number of transactions completed and intended (the latter being just the product of number of clients and number of transactions per client); these will be equal unless the run failed before completion. (In -T mode, only the actual number of transactions is printed.) The last two lines report the number of transactions per second, figured with and without counting the time to start database sessions.

The default TPC-B-like transaction test requires specific tables to be set up beforehand. pgbench should be invoked with the -i (initialize) option to create and populate these tables. (When you are testing a custom script, you don't need this step, but will instead need to do whatever setup your test needs.) Initialization looks like:

pgbench -i [other-options] dbname

where dbname is the name of the already-created database to test in. (You may also need -h, -p, and/or -U options to specify how to connect to the database server.)

Caution

pgbench -i creates four tables pgbench_accounts, pgbench_branches, pgbench_history, and pgbench_tellers, destroying any existing tables of these names. Be very careful to use another database if you have tables having these names!
At the default “scale factor” of 1, the tables initially contain this many rows:

<table>
<thead>
<tr>
<th>table</th>
<th># of rows</th>
</tr>
</thead>
<tbody>
<tr>
<td>pgbench_branches</td>
<td>1</td>
</tr>
<tr>
<td>pgbench_tellers</td>
<td>10</td>
</tr>
<tr>
<td>pgbench_accounts</td>
<td>100000</td>
</tr>
<tr>
<td>pgbench_history</td>
<td>0</td>
</tr>
</tbody>
</table>

You can (and, for most purposes, probably should) increase the number of rows by using the \(-s\) (scale factor) option. The \(-f\) (fillfactor) option might also be used at this point.

Once you have done the necessary setup, you can run your benchmark with a command that doesn’t include \(-i\), that is

```
pgbench [ options ] dbname
```

In nearly all cases, you’ll need some options to make a useful test. The most important options are \(-c\) (number of clients), \(-t\) (number of transactions), \(-T\) (time limit), and \(-f\) (specify a custom script file). See below for a full list.

## Options

The following is divided into three subsections. Different options are used during database initialization and while running benchmarks, but some options are useful in both cases.

### Initialization Options

pgbench accepts the following command-line initialization arguments:

- \(-i\)
  --initialize

    Required to invoke initialization mode.

- \(-I\) init_steps
  --init-steps=init_steps

    Perform just a selected set of the normal initialization steps. \(init\_steps\) specifies the initialization steps to be performed, using one character per step. Each step is invoked in the specified order. The default is \(dtgvp\). The available steps are:

    **d** (Drop)
    Drop any existing pgbench tables.

    **t** (create Tables)
    Create the tables used by the standard pgbench scenario, namely \(pgbench\_accounts\), \(pgbench\_branches\), \(pgbench\_history\), and \(pgbench\_tellers\).

    **g** (Generate data)
    Generate data and load it into the standard tables, replacing any data already present.

    **v** (Vacuum)
    Invoke VACUUM on the standard tables.
p (create Primary keys)

Create primary key indexes on the standard tables.

f (create Foreign keys)

Create foreign key constraints between the standard tables. (Note that this step is not performed by default.)

-F fillfactor
--fillfactor=fillfactor

Create the pgbench_accounts, pgbench_tellers and pgbench_branches tables with the given fillfactor. Default is 100.

-n
--no-vacuum

Perform no vacuuming during initialization. (This option suppresses the v initialization step, even if it was specified in -I.)

-q
--quiet

Switch logging to quiet mode, producing only one progress message per 5 seconds. The default logging prints one message each 100000 rows, which often outputs many lines per second (especially on good hardware).

-s scale_factor
--scale=scale_factor

Multiply the number of rows generated by the scale factor. For example, -s 100 will create 10,000,000 rows in the pgbench_accounts table. Default is 1. When the scale is 20,000 or larger, the columns used to hold account identifiers (aid columns) will switch to using larger integers (bigint), in order to be big enough to hold the range of account identifiers.

--foreign-keys

Create foreign key constraints between the standard tables. (This option adds the f step to the initialization step sequence, if it is not already present.)

--index-tablespace=index_tablespace

Create indexes in the specified tablespace, rather than the default tablespace.

--tablespace=tablespace

Create tables in the specified tablespace, rather than the default tablespace.

--unlogged-tables

Create all tables as unlogged tables, rather than permanent tables.

**Benchmarking Options**

pgbench accepts the following command-line benchmarking arguments:

-b scriptname[@weight]
--builtin=scriptname[@weight]

Add the specified built-in script to the list of executed scripts. An optional integer weight after @ allows to adjust the probability of drawing the script. If not specified, it is set to 1. Available built-
in scripts are: tpch-like, simple-update and select-only. Unambiguous prefixes of built-in names are accepted. With special name list, show the list of built-in scripts and exit immediately.

 `-c clients`
 `-c client=clients`

 Number of clients simulated, that is, number of concurrent database sessions. Default is 1.

 `-C`
 `-C --connect`

 Establish a new connection for each transaction, rather than doing it just once per client session. This is useful to measure the connection overhead.

 `-d`
 `-d --debug`

 Print debugging output.

 `-D varname=value`
 `-D --define=varname=value`

 Define a variable for use by a custom script (see below). Multiple `-D` options are allowed.

 `-f filename[@weight]`
 `-f --file=filename[@weight]`

 Add a transaction script read from `filename` to the list of executed scripts. An optional integer weight after `@` allows to adjust the probability of drawing the test. See below for details.

 `-j threads`
 `-j --jobs=threads`

 Number of worker threads within pgbench. Using more than one thread can be helpful on multi-CPU machines. Clients are distributed as evenly as possible among available threads. Default is 1.

 `-l`
 `-l --log`

 Write information about each transaction to a log file. See below for details.

 `-L limit`
 `-L --latency-limit=limit`

 Transaction which last more than `limit` milliseconds are counted and reported separately, as `late`.

 When throttling is used (`--rate=...`), transactions that lag behind schedule by more than `limit` ms, and thus have no hope of meeting the latency limit, are not sent to the server at all. They are counted and reported separately as `skipped`.

 `-M querymode`
 `-M --protocol=querymode`

 Protocol to use for submitting queries to the server:

 • `simple`: use simple query protocol.
• extended: use extended query protocol.

• prepared: use extended query protocol with prepared statements.
The default is simple query protocol. (See Chapter 53 for more information.)

-n
--no-vacuum

Perform no vacuuming before running the test. This option is necessary if you are running a custom test scenario that does not include the standard tables pgbench_accounts, pgbench_branches, pgbench_history, and pgbench_tellers.

-N
--skip-some-updates


-P sec
--progress=sec

Show progress report every sec seconds. The report includes the time since the beginning of the run, the TPS since the last report, and the transaction latency average and standard deviation since the last report. Under throttling (-R), the latency is computed with respect to the transaction scheduled start time, not the actual transaction beginning time, thus it also includes the average schedule lag time.

-r
--report-latencies

Report the average per-statement latency (execution time from the perspective of the client) of each command after the benchmark finishes. See below for details.

-R rate
--rate=rate

Execute transactions targeting the specified rate instead of running as fast as possible (the default). The rate is given in transactions per second. If the targeted rate is above the maximum possible rate, the rate limit won't impact the results.

The rate is targeted by starting transactions along a Poisson-distributed schedule time line. The expected start time schedule moves forward based on when the client first started, not when the previous transaction ended. That approach means that when transactions go past their original scheduled end time, it is possible for later ones to catch up again.

When throttling is active, the transaction latency reported at the end of the run is calculated from the scheduled start times, so it includes the time each transaction had to wait for the previous transaction to finish. The wait time is called the schedule lag time, and its average and maximum are also reported separately. The transaction latency with respect to the actual transaction start time, i.e. the time spent executing the transaction in the database, can be computed by subtracting the schedule lag time from the reported latency.

If --latency-limit is used together with --rate, a transaction can lag behind so much that it is already over the latency limit when the previous transaction ends, because the latency is calculated from the scheduled start time. Such transactions are not sent to the server, but are skipped altogether and counted separately.

A high schedule lag time is an indication that the system cannot process transactions at the specified rate, with the chosen number of clients and threads. When the average transaction execution time is longer than the scheduled interval between each transaction, each successive transaction
will fall further behind, and the schedule lag time will keep increasing the longer the test run is. When that happens, you will have to reduce the specified transaction rate.

-s scale \_factor
--scale=scale \_factor

Report the specified scale factor in pgbench’s output. With the built-in tests, this is not necessary; the correct scale factor will be detected by counting the number of rows in the pgbench \_branches table. However, when testing only custom benchmarks (\-f option), the scale factor will be reported as 1 unless this option is used.

-S
--select-only

Run built-in select-only script. Shorthand for \-b select-only.

-t transactions
--transactions=transactions

Number of transactions each client runs. Default is 10.

-T seconds
--time=seconds

Run the test for this many seconds, rather than a fixed number of transactions per client. \-t and \-T are mutually exclusive.

-v
--vacuum-all

Vacuum all four standard tables before running the test. With neither \-n nor \-v, pgbench will vacuum the pgbench \_tellers and pgbench \_branches tables, and will truncate pgbench \_history.

--aggregate-interval=seconds

Length of aggregation interval (in seconds). May be used only with \-I option. With this option, the log contains per-interval summary data, as described below.

--log-prefix=prefix

Set the filename prefix for the log files created by \-log. The default is pgbench \_log.

--progress-timestamp

When showing progress (option \-P), use a timestamp (Unix epoch) instead of the number of seconds since the beginning of the run. The unit is in seconds, with millisecond precision after the dot. This helps compare logs generated by various tools.

--random-seed=SEED

Set random generator seed. Seeds the system random number generator, which then produces a sequence of initial generator states, one for each thread. Values for SEED may be: time (the default, the seed is based on the current time), rand (use a strong random source, failing if none is available), or an unsigned decimal integer value. The random generator is invoked explicitly from a pgbench script (random... functions) or implicitly (for instance option \-rate uses it to schedule transactions). When explicitly set, the value used for seeding is shown on the terminal. Any value allowed for SEED may also be provided through the environment variable PGBENCH \_RANDOM \_SEED. To ensure that the provided seed impacts all possible uses, put this option first or use the environment variable.
Setting the seed explicitly allows to reproduce a *pgbench* run exactly, as far as random numbers are concerned. As the random state is managed per thread, this means the exact same *pgbench* run for an identical invocation if there is one client per thread and there are no external or data dependencies. From a statistical viewpoint reproducing runs exactly is a bad idea because it can hide the performance variability or improve performance unduly, e.g. by hitting the same pages as a previous run. However, it may also be of great help for debugging, for instance re-running a tricky case which leads to an error. Use wisely.

```
--sampling-rate=rate
```

Sampling rate, used when writing data into the log, to reduce the amount of log generated. If this option is given, only the specified fraction of transactions are logged. 1.0 means all transactions will be logged, 0.05 means only 5% of the transactions will be logged.

Remember to take the sampling rate into account when processing the log file. For example, when computing TPS values, you need to multiply the numbers accordingly (e.g. with 0.01 sample rate, you'll only get 1/100 of the actual TPS).

**Common Options**

*pgbench* accepts the following command-line common arguments:

```
-h hostname
--host=hostname
```

The database server's host name

```
-p port
--port=port
```

The database server's port number

```
-U login
--username=login
```

The user name to connect as

```
-V
--version
```

Print the *pgbench* version and exit.

```
-?
--help
```

Show help about *pgbench* command line arguments, and exit.

**Notes**

**What is the “Transaction” Actually Performed in pgbench?**

*pgbench* executes test scripts chosen randomly from a specified list. They include built-in scripts with `-b` and user-provided custom scripts with `-f`. Each script may be given a relative weight specified after a `@` so as to change its drawing probability. The default weight is 1. Scripts with a weight of 0 are ignored.

The default built-in transaction script (also invoked with `-b tpcb-like`) issues seven commands per transaction over randomly chosen `aid`, `tid`, `bid` and `balance`. The scenario is inspired by the TPC-B benchmark, but is not actually TPC-B, hence the name.
BEGIN;

UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;

SELECT abalance FROM pgbench_accounts WHERE aid = :aid;

UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;

UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;

INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);

END;

If you select the simple-update built-in (also \texttt{-N}), steps 4 and 5 aren’t included in the transaction. This will avoid update contention on these tables, but it makes the test case even less like TPC-B.

If you select the select-only built-in (also \texttt{-S}), only the \texttt{SELECT} is issued.

### Custom Scripts

pgbench has support for running custom benchmark scenarios by replacing the default transaction script (described above) with a transaction script read from a file (\texttt{-f} option). In this case a “transaction” counts as one execution of a script file.

A script file contains one or more SQL commands terminated by semicolons. Empty lines and lines beginning with \texttt{--} are ignored. Script files can also contain “meta commands”, which are interpreted by pgbench itself, as described below.

**Note**

Before PostgreSQL 9.6, SQL commands in script files were terminated by newlines, and so they could not be continued across lines. Now a semicolon is required to separate consecutive SQL commands (though a SQL command does not need one if it is followed by a meta command). If you need to create a script file that works with both old and new versions of pgbench, be sure to write each SQL command on a single line ending with a semicolon.

There is a simple variable-substitution facility for script files. Variable names must consist of letters (including non-Latin letters), digits, and underscores. Variables can be set by the command-line \texttt{-D} option, explained above, or by the meta commands explained below. In addition to any variables preset by \texttt{-D} command-line options, there are a few variables that are preset automatically, listed in Table 241. A value specified for these variables using \texttt{-D} takes precedence over the automatic presets. Once set, a variable’s value can be inserted into a SQL command by writing \texttt{:variablename}. When running more than one client session, each session has its own set of variables.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>client_id</td>
<td>unique number identifying the client session (starts from zero)</td>
</tr>
<tr>
<td>default_seed</td>
<td>seed used in hash functions by default</td>
</tr>
<tr>
<td>Variable</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>-------------------------------------------------</td>
</tr>
<tr>
<td>random_seed</td>
<td>random generator seed (unless overwritten with -D)</td>
</tr>
<tr>
<td>scale</td>
<td>current scale factor</td>
</tr>
</tbody>
</table>

Script file meta commands begin with a backslash (\) and normally extend to the end of the line, although they can be continued to additional lines by writing backslash-return. Arguments to a meta command are separated by white space. These meta commands are supported:

```
\if expression
\elif expression
\else
\endif
```

This group of commands implements nestable conditional blocks, similarly to psql’s `if expression`. Conditional expressions are identical to those with `set`, with non-zero values interpreted as true.

```
\set varname expression
```

Sets variable `varname` to a value calculated from `expression`. The expression may contain the NULL constant, Boolean constants `TRUE` and `FALSE`, integer constants such as `5432`, double constants such as `3.14159`, references to variables `:variablename`, operators with their usual SQL precedence and associativity, function calls, SQL `CASE` generic conditional expressions and parentheses.

Functions and most operators return NULL on NULL input.

For conditional purposes, non zero numerical values are `TRUE`, zero numerical values and NULL are `FALSE`.

When no final `ELSE` clause is provided to a `CASE`, the default value is NULL.

Examples:

```
\set ntellers 10 * :scale
\set aid (1021 * random(1, 100000 * :scale)) % (100000 * :scale) + 1
\set divx CASE WHEN :x <> 0 THEN :y/:x ELSE NULL END
```

```
\sleep number [ us | ms | s ]
```

Causes script execution to sleep for the specified duration in microseconds (`us`), milliseconds (`ms`) or seconds (`s`). If the unit is omitted then seconds are the default. `number` can be either an integer constant or a `:variablename` reference to a variable having an integer value.

Example:

```
\sleep 10 ms
```

```
\setshell varname command [ argument ... ]
```

Sets variable `varname` to the result of the shell command `command` with the given `argument(s)`. The command must return an integer value through its standard output.

`command` and each `argument` can be either a text constant or a `:variablename` reference to a variable. If you want to use an `argument` starting with a colon, write an additional colon at the beginning of `argument`. 1850
Example:

\setshell variable_to_be_assigned command

literal_argument :variable ::literal_starting_with_colon

\shell command [ argument ... ]

Same as \setshell, but the result of the command is discarded.

Example:

\shell command

literal_argument :variable ::literal_starting_with_colon

### Built-In Operators

The arithmetic, bitwise, comparison and logical operators listed in Table 242 are built into pgbench and may be used in expressions appearing in \set.

#### Table 242. pgbench Operators by increasing precedence

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>OR</td>
<td>logical or</td>
<td>5 or 0</td>
<td>TRUE</td>
</tr>
<tr>
<td>AND</td>
<td>logical and</td>
<td>3 and 0</td>
<td>FALSE</td>
</tr>
<tr>
<td>NOT</td>
<td>logical not</td>
<td>not false</td>
<td>TRUE</td>
</tr>
<tr>
<td>IS [NOT] (NULL</td>
<td>TRUE</td>
<td>FALSE)</td>
<td>value tests</td>
</tr>
<tr>
<td>ISNULL</td>
<td>NOTNULL</td>
<td>null tests</td>
<td>1 notnull</td>
</tr>
<tr>
<td>=</td>
<td>is equal</td>
<td>5 = 4</td>
<td>FALSE</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>is not equal</td>
<td>5 &lt;&gt; 4</td>
<td>TRUE</td>
</tr>
<tr>
<td>!=</td>
<td>is not equal</td>
<td>5 != 5</td>
<td>FALSE</td>
</tr>
<tr>
<td>&lt;</td>
<td>lower than</td>
<td>5 &lt; 4</td>
<td>FALSE</td>
</tr>
<tr>
<td>&lt;=</td>
<td>lower or equal</td>
<td>5 &lt;= 4</td>
<td>FALSE</td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
<td>5 &gt; 4</td>
<td>TRUE</td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater or equal</td>
<td>5 &gt;= 4</td>
<td>TRUE</td>
</tr>
<tr>
<td></td>
<td></td>
<td>integer bitwise OR</td>
<td>1</td>
</tr>
<tr>
<td>#</td>
<td>integer bitwise XOR</td>
<td>1 # 3</td>
<td>2</td>
</tr>
<tr>
<td>&amp;</td>
<td>integer bitwise AND</td>
<td>1 &amp; 3</td>
<td>1</td>
</tr>
<tr>
<td>~</td>
<td>integer bitwise NOT</td>
<td>~ 1</td>
<td>-2</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>integer bitwise shift left</td>
<td>1 &lt;&lt; 2</td>
<td>4</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>integer bitwise shift right</td>
<td>8 &gt;&gt; 2</td>
<td>2</td>
</tr>
<tr>
<td>+</td>
<td>addition</td>
<td>5 + 4</td>
<td>9</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>3 - 2.0</td>
<td>1.0</td>
</tr>
<tr>
<td>*</td>
<td>multiplication</td>
<td>5 * 4</td>
<td>20</td>
</tr>
<tr>
<td>/</td>
<td>division (integer truncates the results)</td>
<td>5 / 3</td>
<td>1</td>
</tr>
<tr>
<td>%</td>
<td>modulo</td>
<td>3 % 2</td>
<td>1</td>
</tr>
</tbody>
</table>
### Built-In Functions

The functions listed in Table 243 are built into pgbench and may be used in expressions appearing in `\set`.

#### Table 243. pgbench Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>abs(a)</td>
<td>same as a</td>
<td>absolute value</td>
<td>abs(-17)</td>
<td>17</td>
</tr>
<tr>
<td>debug(a)</td>
<td>same as a</td>
<td>print a to stderr, and return a</td>
<td>debug(5432.1)</td>
<td>5432.1</td>
</tr>
<tr>
<td>double(i)</td>
<td>double</td>
<td>cast to double</td>
<td>double(5432)</td>
<td>5432.0</td>
</tr>
<tr>
<td>exp(x)</td>
<td>double</td>
<td>exponential</td>
<td>exp(1.0)</td>
<td>2.7182818284590459045</td>
</tr>
<tr>
<td>greatest(a [, ...])</td>
<td>double if any a is largest value among arguments</td>
<td>greatest(5, 4, 3, 2)</td>
<td>5</td>
<td></td>
</tr>
<tr>
<td>hash(a [seed])</td>
<td>integer</td>
<td>alias for hash_mur-mur2()</td>
<td>hash(10, 5432)</td>
<td>-5817877081768721676</td>
</tr>
<tr>
<td>hash_fnv1a(a [seed])</td>
<td>integer</td>
<td>FNV-1a hash¹</td>
<td>hash_fnv1a(10, 5432)</td>
<td>-7793829335365542153</td>
</tr>
<tr>
<td>hash_mur-mur2(a [seed])</td>
<td>integer</td>
<td>MurmurHash2 hash²</td>
<td>hash_mur-mur2(10, 5432)</td>
<td>-5817877081768721676</td>
</tr>
<tr>
<td>int(x)</td>
<td>integer</td>
<td>cast to int</td>
<td>int(5.4, 3.8)</td>
<td>+ 9</td>
</tr>
<tr>
<td>least(a [, ...])</td>
<td>double if any a is smallest value among arguments</td>
<td>least(5, 4, 2.1)</td>
<td>2.1</td>
<td></td>
</tr>
<tr>
<td>ln(x)</td>
<td>double</td>
<td>natural logarithm</td>
<td>ln(2.718281828459045)</td>
<td>0.99959045</td>
</tr>
<tr>
<td>mod(i, j)</td>
<td>integer</td>
<td>modulo</td>
<td>mod(54, 32)</td>
<td>22</td>
</tr>
<tr>
<td>pi()</td>
<td>double</td>
<td>value of the constant PI</td>
<td>pi()</td>
<td>3.14159265358979323846</td>
</tr>
<tr>
<td>pow(x, y), power(x, y)</td>
<td>double</td>
<td>exponentiation</td>
<td>pow(2.0, 10), 1024.0</td>
<td></td>
</tr>
<tr>
<td>random(lb, ub)</td>
<td>integer</td>
<td>uniformly-distributed random integer in [lb, ub]</td>
<td>random(1, 10)</td>
<td>an integer between 1 and 10</td>
</tr>
<tr>
<td>random_exponential(lb, ub, parameter)</td>
<td>integer</td>
<td>exponentially-distributed random integer in [lb, ub]</td>
<td>random_exponential(1, 10, 3.0)</td>
<td>an integer between 1 and 10</td>
</tr>
<tr>
<td>random_gauss-</td>
<td>integer</td>
<td>Gaussian-distributed random integer</td>
<td>random_gauss-</td>
<td>an integer between 1 and 10</td>
</tr>
</tbody>
</table>

¹ [https://en.wikipedia.org/wiki/Fowler%E2%80%93Noll%E2%80%93Vo_hash_function](https://en.wikipedia.org/wiki/Fowler%E2%80%93Noll%E2%80%93Vo_hash_function)

² [https://en.wikipedia.org/wiki/MurmurHash](https://en.wikipedia.org/wiki/MurmurHash)
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ian(lb, ub, parameter)</code></td>
<td>integer</td>
<td>Zipfian-distributed random integer in <code>[lb, ub]</code>, see below</td>
<td><code>ian(1, 10, 1)</code></td>
<td></td>
</tr>
<tr>
<td><code>random_zipfian(lb, ub, parameter)</code></td>
<td>integer</td>
<td>Approximated bounded Zipfian distribution. For <code>parameter</code> in <code>(0, 1)</code>, an approximated algorithm is taken from &quot;Quickly Generating Billion-Record Synthetic Databases&quot;, Jim Gray et al, SIGMOD 1994. For <code>parameter</code> in <code>(1, 1000)</code>, a rejection method is used, based on &quot;Non-Uniform Random Variate Generation&quot;, Luc Devroye, p. 550-551, Springer 1986. The distribution is not defined when the parameter's value is 1.0. The drawing performance is poor for parameter values close and above 1.0 and on a small range. <code>parameter</code> defines how skewed the distribution is. The larger the <code>parameter</code>, the more frequently values to the beginning of the interval are drawn. The closer to 0 <code>parameter</code> is, the flatter (more uniform) the access distribution.</td>
<td><code>random_zipfian(1, 10, 1)</code> below</td>
<td></td>
</tr>
</tbody>
</table>

The `random` function generates values using a uniform distribution, that is all the values are drawn within the specified range with equal probability. The `random_exponential`, `random_gaussian` and `random_zipfian` functions require an additional double parameter which determines the precise shape of the distribution.

- For an exponential distribution, `parameter` controls the distribution by truncating a quickly-decreasing exponential distribution at `parameter`, and then projecting onto integers between the bounds. To be precise, with

\[
f(x) = \exp(-parameter \times (x - min) / (max - min + 1)) / (1 - \exp(-parameter))
\]

Then value `i` between `min` and `max` inclusive is drawn with probability: \( f(i) - f(i + 1) \).

Intuitively, the larger the `parameter`, the more frequently values close to `min` are accessed, and the less frequently values close to `max` are accessed. The closer to 0 `parameter` is, the flatter (more uniform) the access distribution. A crude approximation of the distribution is that the most frequent 1% values in the range, close to `min`, are drawn `parameter`% of the time. The `parameter` value must be strictly positive.

- For a Gaussian distribution, the interval is mapped onto a standard normal distribution (the classical bell-shaped Gaussian curve) truncated at `-parameter` on the left and `+parameter` on the right. Values in the middle of the interval are more likely to be drawn. To be precise, if \( \Phi(x) \) is the cumulative distribution function of the standard normal distribution, with mean \( mu \) defined as \( (max + min) / 2.0 \), with

\[
f(x) = \Phi(2.0 \times parameter \times (x - mu) / (max - min + 1)) / (2.0 \times \Phi(parameter) - 1)
\]

then value `i` between `min` and `max` inclusive is drawn with probability: \( f(i + 0.5) - f(i - 0.5) \). Intuitively, the larger the `parameter`, the more frequently values close to the middle of the interval are drawn, and the less frequently values close to the `min` and `max` bounds. About 67% of values are drawn from the middle 1.0 / `parameter`, that is a relative 0.5 / `parameter` around the mean, and 95% in the middle 2.0 / `parameter`, that is a relative 1.0 / `parameter` around the mean; for instance, if `parameter` is 4.0, 67% of values are drawn from the middle quarter (1.0 / 4.0) of the interval (i.e. from 3.0 / 8.0 to 5.0 / 8.0) and 95% from the middle half (2.0 / 4.0) of the interval (second and third quartiles). The minimum `parameter` is 2.0 for performance of the Box-Muller transform.

- `random_zipfian` generates an approximated bounded Zipfian distribution. For `parameter` in `(0, 1)`, an approximated algorithm is taken from "Quickly Generating Billion-Record Synthetic Databases", Jim Gray et al, SIGMOD 1994. For `parameter` in `(1, 1000)`, a rejection method is used, based on "Non-Uniform Random Variate Generation", Luc Devroye, p. 550-551, Springer 1986. The distribution is not defined when the parameter's value is 1.0. The drawing performance is poor for parameter values close and above 1.0 and on a small range.

`parameter` defines how skewed the distribution is. The larger the `parameter`, the more frequently values to the beginning of the interval are drawn. The closer to 0 `parameter` is, the flatter (more uniform) the access distribution.
Hash functions `hash`, `hash_murmur2` and `hash_fnv1a` accept an input value and an optional seed parameter. In case the seed isn’t provided the value of `:default_seed` is used, which is initialized randomly unless set by the command-line `-D` option. Hash functions can be used to scatter the distribution of random functions such as `random_zipfian` or `random_exponential`. For instance, the following pgbench script simulates possible real world workload typical for social media and blogging platforms where few accounts generate excessive load:

```sql
\set r random_zipfian(0, 100000000, 1.07)
\set k abs(hash(:r)) % 1000000
```

In some cases several distinct distributions are needed which don’t correlate with each other and this is when implicit seed parameter comes in handy:

```sql
\set k1 abs(hash(:r, :default_seed + 123)) % 1000000
\set k2 abs(hash(:r, :default_seed + 321)) % 1000000
```

As an example, the full definition of the built-in TPC-B-like transaction is:

```sql
\set aid random(1, 100000 * :scale)
\set bid random(1, 1 * :scale)
\set tid random(1, 10 * :scale)
\set delta random(-5000, 5000)
BEGIN;
UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;
SELECT abalance FROM pgbench_accounts WHERE aid = :aid;
UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;
UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;
INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);
END;
```

This script allows each iteration of the transaction to reference different, randomly-chosen rows. (This example also shows why it’s important for each client session to have its own variables — otherwise they’d not be independently touching different rows.)

### Per-Transaction Logging

With the `-l` option (but without the `--aggregate-interval` option), pgbench writes information about each transaction to a log file. The log file will be named `prefix.nnn`, where `prefix` defaults to `pgbench_log`, and `nnn` is the PID of the pgbench process. The prefix can be changed by using the `--log-prefix` option. If the `-j` option is 2 or higher, so that there are multiple worker threads, each will have its own log file. The first worker will use the same name for its log file as in the standard single worker case. The additional log files for the other workers will be named `prefix.nnn.mmm`, where `mmm` is a sequential number for each worker starting with 1.

The format of the log is:

```plaintext
client_id transaction_no time script_no time_epoch time_us
[ schedule_lag ]
```

where `client_id` indicates which client session ran the transaction, `transaction_no` counts how many transactions have been run by that session, `time` is the total elapsed transaction time in microseconds, `script_no` identifies which script file was used (useful when multiple scripts were
specified with -f or -b), and time_epoch/time_us are a Unix-epoch time stamp and an offset in microseconds (suitable for creating an ISO 8601 time stamp with fractional seconds) showing when the transaction completed. The schedule_lag field is the difference between the transaction’s scheduled start time, and the time it actually started, in microseconds. It is only present when the --rate option is used. When both --rate and --latency-limit are used, the time for a skipped transaction will be reported as skipped.

Here is a snippet of a log file generated in a single-client run:

```
0 199 2241 0 1175850568 995598
0 200 2465 0 1175850568 998079
0 201 2513 0 1175850569 608
0 202 2038 0 1175850569 2663
```

Another example with --rate=100 and --latency-limit=5 (note the additional schedule_lag column):

```
0 81 4621 0 1412881037 912698 3005
0 82 6173 0 1412881037 914578 4304
0 83 skipped 0 1412881037 914578 5217
0 83 skipped 0 1412881037 914578 5099
0 83 4722 0 1412881037 916203 3108
0 84 4142 0 1412881037 918023 2333
0 85 2465 0 1412881037 919759 740
```

In this example, transaction 82 was late, because its latency (6.173 ms) was over the 5 ms limit. The next two transactions were skipped, because they were already late before they were even started.

When running a long test on hardware that can handle a lot of transactions, the log files can become very large. The --sampling-rate option can be used to log only a random sample of transactions.

**Aggregated Logging**

With the --aggregate-interval option, a different format is used for the log files:

```
interval_start num_transactions sum_latency sum_latency_2 min_latency max_latency
[ sum_lag sum_lag_2 min_lag max_lag [ skipped ] ]
```

where interval_start is the start of the interval (as a Unix epoch time stamp), num_transactions is the number of transactions within the interval, sum_latency is the sum of the transaction latencies within the interval, sum_latency_2 is the sum of squares of the transaction latencies within the interval, min_latency is the minimum latency within the interval, and max_latency is the maximum latency within the interval. The next fields, sum_lag, sum_lag_2, min_lag, and max_lag, are only present if the --rate option is used. They provide statistics about the time each transaction had to wait for the previous one to finish, i.e. the difference between each transaction's scheduled start time and the time it actually started. The very last field, skipped, is only present if the --latency-limit option is used, too. It counts the number of transactions skipped because they would have started too late. Each transaction is counted in the interval when it was committed.

Here is some example output:

```
1345828501 5601 1542744 483552416 61 2573
1345828503 7884 1979812 565806736 60 1479
1345828505 7208 1979422 573489941 236 1411
```

1855
Notice that while the plain (unaggregated) log file shows which script was used for each transaction, the aggregated log does not. Therefore if you need per-script data, you need to aggregate the data on your own.

**Per-Statement Latencies**

With the `--r` option, pgbench collects the elapsed transaction time of each statement executed by every client. It then reports an average of those values, referred to as the latency for each statement, after the benchmark has finished.

For the default script, the output will look similar to this:

```
starting vacuum...end.
transaction type: <builtin: TPC-B (sort of)>
scaling factor: 1
query mode: simple
number of clients: 10
number of threads: 1
number of transactions per client: 1000
number of transactions actually processed: 10000/10000
latency average = 15.844 ms
latency stddev = 2.715 ms
tps = 618.764555 (including connections establishing)
tps = 622.977698 (excluding connections establishing)
statement latencies in milliseconds:
  0.002  \set aid random(1, 100000 * :scale)
  0.005  \set bid random(1, 1 * :scale)
  0.002  \set tid random(1, 10 * :scale)
  0.001  \set delta random(-5000, 5000)
  0.326  BEGIN;
  0.603  UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;
  0.454  SELECT abalance FROM pgbench_accounts WHERE aid = :aid;
  5.528  UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;
  7.335  UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;
  0.371  INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);
  1.212  END;
```

If multiple script files are specified, the averages are reported separately for each script file.

Note that collecting the additional timing information needed for per-statement latency computation adds some overhead. This will slow average execution speed and lower the computed TPS. The amount of slowdown varies significantly depending on platform and hardware. Comparing average TPS values with and without latency reporting enabled is a good way to measure if the timing overhead is significant.

**Good Practices**

It is very easy to use pgbench to produce completely meaningless numbers. Here are some guidelines to help you get useful results.

In the first place, *never* believe any test that runs for only a few seconds. Use the `--t` or `--T` option to make the run last at least a few minutes, so as to average out noise. In some cases you could need...
hours to get numbers that are reproducible. It's a good idea to try the test run a few times, to find out if your numbers are reproducible or not.

For the default TPC-B-like test scenario, the initialization scale factor (-s) should be at least as large as the largest number of clients you intend to test (-c); else you'll mostly be measuring update contention. There are only -s rows in the pgbench_branches table, and every transaction wants to update one of them, so -c values in excess of -s will undoubtedly result in lots of transactions blocked waiting for other transactions.

The default test scenario is also quite sensitive to how long it's been since the tables were initialized: accumulation of dead rows and dead space in the tables changes the results. To understand the results you must keep track of the total number of updates and when vacuuming happens. If autovacuum is enabled it can result in unpredictable changes in measured performance.

A limitation of pgbench is that it can itself become the bottleneck when trying to test a large number of client sessions. This can be alleviated by running pgbench on a different machine from the database server, although low network latency will be essential. It might even be useful to run several pgbench instances concurrently, on several client machines, against the same database server.

Security

If untrusted users have access to a database that has not adopted a secure schema usage pattern, do not run pgbench in that database. pgbench uses unqualified names and does not manipulate the search path.
pg_config

pg_config — retrieve information about the installed version of PostgreSQL

Synopsis

pg_config [option...]

Description

The pg_config utility prints configuration parameters of the currently installed version of PostgreSQL. It is intended, for example, to be used by software packages that want to interface to PostgreSQL to facilitate finding the required header files and libraries.

Options

To use pg_config, supply one or more of the following options:

--bindir

Print the location of user executables. Use this, for example, to find the psql program. This is normally also the location where the pg_config program resides.

--docdir

Print the location of documentation files.

--htmldir

Print the location of HTML documentation files.

--includedir

Print the location of C header files of the client interfaces.

--pkgincludedir

Print the location of other C header files.

--includedir-server

Print the location of C header files for server programming.

--libdir

Print the location of object code libraries.

--pkglibdir

Print the location of dynamically loadable modules, or where the server would search for them. (Other architecture-dependent data files might also be installed in this directory.)

--localedir

Print the location of locale support files. (This will be an empty string if locale support was not configured when PostgreSQL was built.)
--mandir
Print the location of manual pages.

--sharedir
Print the location of architecture-independent support files.

--sysconfdir
Print the location of system-wide configuration files.

--pgxs
Print the location of extension makefiles.

--configure
Print the options that were given to the configure script when PostgreSQL was configured for building. This can be used to reproduce the identical configuration, or to find out with what options a binary package was built. (Note however that binary packages often contain vendor-specific custom patches.) See also the examples below.

--cc
Print the value of the CC variable that was used for building PostgreSQL. This shows the C compiler used.

--cppflags
Print the value of the CPPFLAGS variable that was used for building PostgreSQL. This shows C compiler switches needed at preprocessing time (typically, -I switches).

--cflags
Print the value of the CFLAGS variable that was used for building PostgreSQL. This shows C compiler switches.

--cflags_sl
Print the value of the CFLAGS_SL variable that was used for building PostgreSQL. This shows extra C compiler switches used for building shared libraries.

--ldflags
Print the value of the LDFLAGS variable that was used for building PostgreSQL. This shows linker switches.

--ldflags_ex
Print the value of the LDFLAGS_EX variable that was used for building PostgreSQL. This shows linker switches used for building executables only.

--ldflags_sl
Print the value of the LDFLAGS_SL variable that was used for building PostgreSQL. This shows linker switches used for building shared libraries only.

--libs
Print the value of the LIBS variable that was used for building PostgreSQL. This normally contains -l switches for external libraries linked into PostgreSQL.
pg_config

--version

Print the version of PostgreSQL.

-?
--help

Show help about pg_config command line arguments, and exit.

If more than one option is given, the information is printed in that order, one item per line. If no options are given, all available information is printed, with labels.

Notes

The options --docdir, --pkgincludedir, --localedir, --mandir, --sharedir, --sysconfdir, --cc, --cppflags, --cflags, --cflags_sl, --ldflags, --ldflags_sl, and --libs were added in PostgreSQL 8.1. The option --htmldir was added in PostgreSQL 8.4. The option --ldflags_ex was added in PostgreSQL 9.0.

Example

To reproduce the build configuration of the current PostgreSQL installation, run the following command:

`eval ./configure `pg_config --configure`

The output of pg_config --configure contains shell quotation marks so arguments with spaces are represented correctly. Therefore, using eval is required for proper results.
pg_dump

pg_dump — extract a PostgreSQL database into a script file or other archive file

Synopsis

pg_dump [connection-option...] [option...] [dbname]

Description

pg_dump is a utility for backing up a PostgreSQL database. It makes consistent backups even if the
database is being used concurrently. pg_dump does not block other users accessing the database (read-
ers or writers).

pg_dump only dumps a single database. To back up an entire cluster, or to back up global objects that
are common to all databases in a cluster (such as roles and tablespaces), use pg_dumpall.

Dumps can be output in script or archive file formats. Script dumps are plain-text files containing the
SQL commands required to reconstruct the database to the state it was in at the time it was saved.
To restore from such a script, feed it to psql. Script files can be used to reconstruct the database even
on other machines and other architectures; with some modifications, even on other SQL database
products.

The alternative archive file formats must be used with pg_restore to rebuild the database. They allow
pg_restore to be selective about what is restored, or even to reorder the items prior to being restored.
The archive file formats are designed to be portable across architectures.

When used with one of the archive file formats and combined with pg_restore, pg_dump provides a
flexible archival and transfer mechanism. pg_dump can be used to backup an entire database, then
pg_restore can be used to examine the archive and/or select which parts of the database are to be re-
stored. The most flexible output file formats are the “custom” format (-Fc) and the “directory” format
(-Fd). They allow for selection and reordering of all archived items, support parallel restoration, and
are compressed by default. The “directory” format is the only format that supports parallel dumps.

While running pg_dump, one should examine the output for any warnings (printed on standard error),
especially in light of the limitations listed below.

Options

The following command-line options control the content and format of the output.

dbname

Specifies the name of the database to be dumped. If this is not specified, the environment variable
PGDATABASE is used. If that is not set, the user name specified for the connection is used.

-a
--data-only

Dump only the data, not the schema (data definitions). Table data, large objects, and sequence
values are dumped.

This option is similar to, but for historical reasons not identical to, specifying --section=da-
ta.
-b
--blobs
Include large objects in the dump. This is the default behavior except when --schema, --table, or --schema-only is specified. The -b switch is therefore only useful to add large objects to dumps where a specific schema or table has been requested. Note that blobs are considered data and therefore will be included when --data-only is used, but not when --schema-only is.

-B
--no-blobs
Exclude large objects in the dump.
When both -b and -B are given, the behavior is to output large objects, when data is being dumped, see the -b documentation.

-c
--clean
Output commands to clean (drop) database objects prior to outputting the commands for creating them. (Unless --if-exists is also specified, restore might generate some harmless error messages, if any objects were not present in the destination database.)
This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

-C
--create
Begin the output with a command to create the database itself and reconnect to the created database. (With a script of this form, it doesn't matter which database in the destination installation you connect to before running the script.) If --clean is also specified, the script drops and recreates the target database before reconnecting to it.
With --create, the output also includes the database's comment if any, and any configuration variable settings that are specific to this database, that is, any ALTER DATABASE ... SET ... and ALTER ROLE ... IN DATABASE ... SET ... commands that mention this database. Access privileges for the database itself are also dumped, unless --no-acl is specified.
This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

-E encoding
--encoding=encoding
Create the dump in the specified character set encoding. By default, the dump is created in the database encoding. (Another way to get the same result is to set the PGCLIENTENCODING environment variable to the desired dump encoding.)

-f file
--file=file
Send output to the specified file. This parameter can be omitted for file based output formats, in which case the standard output is used. It must be given for the directory output format however, where it specifies the target directory instead of a file. In this case the directory is created by pg_dump and must not exist before.

-F format
--format=format
Selects the format of the output. format can be one of the following:
pg_dump

p
plain

Output a plain-text SQL script file (the default).

c
custom

Output a custom-format archive suitable for input into pg_restore. Together with the directory output format, this is the most flexible output format in that it allows manual selection and reordering of archived items during restore. This format is also compressed by default.

d
directory

Output a directory-format archive suitable for input into pg_restore. This will create a directory with one file for each table and blob being dumped, plus a so-called Table of Contents file describing the dumped objects in a machine-readable format that pg_restore can read. A directory format archive can be manipulated with standard Unix tools; for example, files in an uncompressed archive can be compressed with the gzip tool. This format is compressed by default and also supports parallel dumps.

t
tar

Output a tar-format archive suitable for input into pg_restore. The tar format is compatible with the directory format: extracting a tar-format archive produces a valid directory-format archive. However, the tar format does not support compression. Also, when using tar format the relative order of table data items cannot be changed during restore.

-j njobs
--jobs=njobs

Run the dump in parallel by dumping njobs tables simultaneously. This option reduces the time of the dump but it also increases the load on the database server. You can only use this option with the directory output format because this is the only output format where multiple processes can write their data at the same time.

pg_dump will open njobs + 1 connections to the database, so make sure your max_connections setting is high enough to accommodate all connections.

Requesting exclusive locks on database objects while running a parallel dump could cause the dump to fail. The reason is that the pg_dump master process requests shared locks on the objects that the worker processes are going to dump later in order to make sure that nobody deletes them and makes them go away while the dump is running. If another client then requests an exclusive lock on a table, that lock will not be granted but will be queued waiting for the shared lock of the master process to be released. Consequently any other access to the table will not be granted either and will queue after the exclusive lock request. This includes the worker process trying to dump the table. Without any precautions this would be a classic deadlock situation. To detect this conflict, the pg_dump worker process requests another shared lock using the NOWAIT option. If the worker process is not granted this shared lock, somebody else must have requested an exclusive lock in the meantime and there is no way to continue with the dump, so pg_dump has no choice but to abort the dump.

For a consistent backup, the database server needs to support synchronized snapshots, a feature that was introduced in PostgreSQL 9.2 for primary servers and 10 for standbys. With this feature, database clients can ensure they see the same data set even though they use different connections. pg_dump -j uses multiple database connections; it connects to the database once with the master process and once again for each worker job. Without the synchronized snapshot feature, the different worker jobs wouldn't be guaranteed to see the same data in each connection, which could lead to an inconsistent backup.
If you want to run a parallel dump of a pre-9.2 server, you need to make sure that the database content doesn't change from between the time the master connects to the database until the last worker job has connected to the database. The easiest way to do this is to halt any data modifying processes (DDL and DML) accessing the database before starting the backup. You also need to specify the `--no-synchronized-snapshots` parameter when running `pg_dump -j` against a pre-9.2 PostgreSQL server.

```
-n schema
--schema=schema
```

Dump only schemas matching `schema`; this selects both the schema itself, and all its contained objects. When this option is not specified, all non-system schemas in the target database will be dumped. Multiple schemas can be selected by writing multiple `-n` switches. Also, the `schema` parameter is interpreted as a pattern according to the same rules used by `psql`'s `\d` commands (see Patterns), so multiple schemas can also be selected by writing wildcard characters in the pattern. When using wildcards, be careful to quote the pattern if needed to prevent the shell from expanding the wildcards; see Examples.

**Note**

When `-n` is specified, `pg_dump` makes no attempt to dump any other database objects that the selected schema(s) might depend upon. Therefore, there is no guarantee that the results of a specific-schema dump can be successfully restored by themselves into a clean database.

```
-N schema
--exclude-schema=schema
```

Do not dump any schemas matching the `schema` pattern. The pattern is interpreted according to the same rules as for `-n`. `-N` can be given more than once to exclude schemas matching any of several patterns.

When both `-n` and `-N` are given, the behavior is to dump just the schemas that match at least one `-n` switch but no `-N` switches. If `-N` appears without `-n`, then schemas matching `-N` are excluded from what is otherwise a normal dump.

```
-o
--oids
```

Dump object identifiers (OIDs) as part of the data for every table. Use this option if your application references the OID columns in some way (e.g., in a foreign key constraint). Otherwise, this option should not be used.

```
-O
--no-owner
```

Do not output commands to set ownership of objects to match the original database. By default, `pg_dump` issues `ALTER OWNER` or `SET SESSION AUTHORIZATION` statements to set ownership of created database objects. These statements will fail when the script is run unless it is started by a superuser (or the same user that owns all of the objects in the script). To make a script that can be restored by any user, but will give that user ownership of all the objects, specify `-O`. 

1864
This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

-R
--no-reconnect

This option is obsolete but still accepted for backwards compatibility.

-s
--schema-only

Dump only the object definitions (schema), not data.

This option is the inverse of --data-only. It is similar to, but for historical reasons not identical to, specifying --section=pre-data --section=post-data.

(Do not confuse this with the --schema option, which uses the word “schema” in a different meaning.)

To exclude table data for only a subset of tables in the database, see --exclude-table-data.

-S username
--superuser=username

Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used. (Usually, it's better to leave this out, and instead start the resulting script as superuser.)

-t table
--table=table

Dump only tables with names matching table. For this purpose, “table” includes views, materialized views, sequences, and foreign tables. Multiple tables can be selected by writing multiple -t switches. Also, the table parameter is interpreted as a pattern according to the same rules used by psql's \d commands (see Patterns), so multiple tables can also be selected by writing wildcard characters in the pattern. When using wildcards, be careful to quote the pattern if needed to prevent the shell from expanding the wildcards; see Examples.

The -n and -N switches have no effect when -t is used, because tables selected by -t will be dumped regardless of those switches, and non-table objects will not be dumped.

**Note**

When -t is specified, pg_dump makes no attempt to dump any other database objects that the selected table(s) might depend upon. Therefore, there is no guarantee that the results of a specific-table dump can be successfully restored by themselves into a clean database.

**Note**

The behavior of the -t switch is not entirely upward compatible with pre-8.2 PostgreSQL versions. Formerly, writing -t tab would dump all tables named tab, but now it just dumps whichever one is visible in your default search path. To get the old behavior you can write -t '*.tab'. Also, you must write something like -t sch.tab to select a table in a particular schema, rather than the old location of -n sch -t tab.
-T table
--exclude-table=table

Do not dump any tables matching the table pattern. The pattern is interpreted according to the same rules as for -t. -T can be given more than once to exclude tables matching any of several patterns.

When both -t and -T are given, the behavior is to dump just the tables that match at least one -t switch but no -T switches. If -T appears without -t, then tables matching -T are excluded from what is otherwise a normal dump.

-v
--verbose

Specifies verbose mode. This will cause pg_dump to output detailed object comments and start/stop times to the dump file, and progress messages to standard error.

-V
--version

Print the pg_dump version and exit.

-x
--no-privileges
--no-acl

Prevent dumping of access privileges (grant/revoke commands).

-Z 0..9
--compress=0..9

Specify the compression level to use. Zero means no compression. For the custom archive format, this specifies compression of individual table-data segments, and the default is to compress at a moderate level. For plain text output, setting a nonzero compression level causes the entire output file to be compressed, as though it had been fed through gzip; but the default is not to compress. The tar archive format currently does not support compression at all.

--binary-upgrade

This option is for use by in-place upgrade utilities. Its use for other purposes is not recommended or supported. The behavior of the option may change in future releases without notice.

--column-inserts
--attribute-inserts

Dump data as INSERT commands with explicit column names (INSERT INTO table (col-
umn, ...) VALUES ...). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases. However, since this option generates a separate command for each row, an error in reloading a row causes only that row to be lost rather than the entire table contents.

--disable-dollar-quoting

This option disables the use of dollar quoting for function bodies, and forces them to be quoted using SQL standard string syntax.

--disable-triggers

This option is relevant only when creating a data-only dump. It instructs pg_dump to include commands to temporarily disable triggers on the target tables while the data is reloaded. Use this if you have referential integrity checks or other triggers on the tables that you do not want to invoke during data reload.
Presently, the commands emitted for --disable-triggers must be done as superuser. So, you should also specify a superuser name with -S, or preferably be careful to start the resulting script as a superuser.

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

--enable-row-security

This option is relevant only when dumping the contents of a table which has row security. By default, pg_dump will set row_security to off, to ensure that all data is dumped from the table. If the user does not have sufficient privileges to bypass row security, then an error is thrown. This parameter instructs pg_dump to set row_security to on instead, allowing the user to dump the parts of the contents of the table that they have access to.

Note that if you use this option currently, you probably also want the dump be in INSERT format, as the COPY FROM during restore does not support row security.

--exclude-table-data=table

Do not dump data for any tables matching the table pattern. The pattern is interpreted according to the same rules as for -t. --exclude-table-data can be given more than once to exclude tables matching any of several patterns. This option is useful when you need the definition of a particular table even though you do not need the data in it.

To exclude data for all tables in the database, see --schema-only.

--if-exists

Use conditional commands (i.e. add an IF EXISTS clause) when cleaning database objects. This option is not valid unless --clean is also specified.

--inserts

Dump data as INSERT commands (rather than COPY). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases. However, since this option generates a separate command for each row, an error in reloading a row causes only that row to be lost rather than the entire table contents. Note that the restore might fail altogether if you have rearranged column order. The --column-inserts option is safe against column order changes, though even slower.

--load-via-partition-root

When dumping data for a table partition, make the COPY or INSERT statements target the root of the partitioning hierarchy that contains it, rather than the partition itself. This causes the appropriate partition to be re-determined for each row when the data is loaded. This may be useful when reloading data on a server where rows do not always fall into the same partitions as they did on the original server. That could happen, for example, if the partitioning column is of type text and the two systems have different definitions of the collation used to sort the partitioning column.

It is best not to use parallelism when restoring from an archive made with this option, because pg_restore will not know exactly which partition(s) a given archive data item will load data into. This could result in inefficiency due to lock conflicts between parallel jobs, or perhaps even reload failures due to foreign key constraints being set up before all the relevant data is loaded.

--lock-wait-timeout=timeout

Do not wait forever to acquire shared table locks at the beginning of the dump. Instead fail if unable to lock a table within the specified timeout. The timeout may be specified in any of the formats accepted by SET statement_timeout. (Allowed formats vary depending on the server version you are dumping from, but an integer number of milliseconds is accepted by all versions.)
**--no-comments**
Do not dump comments.

**--no-publications**
Do not dump publications.

**--no-security-labels**
Do not dump security labels.

**--no-subscriptions**
Do not dump subscriptions.

**--no-sync**
By default, `pg_dump` will wait for all files to be written safely to disk. This option causes `pg_dump` to return without waiting, which is faster, but means that a subsequent operating system crash can leave the dump corrupt. Generally, this option is useful for testing but should not be used when dumping data from production installation.

**--no-synchronized-snapshots**
This option allows running `pg_dump -j` against a pre-9.2 server, see the documentation of the `-j` parameter for more details.

**--no-tablespaces**
Do not output commands to select tablespaces. With this option, all objects will be created in whichever tablespace is the default during restore.

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call `pg_restore`.

**--no-unlogged-table-data**
Do not dump the contents of unlogged tables. This option has no effect on whether or not the table definitions (schema) are dumped; it only suppresses dumping the table data. Data in unlogged tables is always excluded when dumping from a standby server.

**--quote-all-identifiers**
Force quoting of all identifiers. This option is recommended when dumping a database from a server whose PostgreSQL major version is different from `pg_dump`'s, or when the output is intended to be loaded into a server of a different major version. By default, `pg_dump` quotes only identifiers that are reserved words in its own major version. This sometimes results in compatibility issues when dealing with servers of other versions that may have slightly different sets of reserved words. Using `--quote-all-identifiers` prevents such issues, at the price of a harder-to-read dump script.

**--section=sectionname**
Only dump the named section. The section name can be `pre-data`, `data`, or `post-data`. This option can be specified more than once to select multiple sections. The default is to dump all sections.

The data section contains actual table data, large-object contents, and sequence values. Post-data items include definitions of indexes, triggers, rules, and constraints other than validated check constraints. Pre-data items include all other data definition items.
--serializable-deferrable

Use a serializable transaction for the dump, to ensure that the snapshot used is consistent with later database states; but do this by waiting for a point in the transaction stream at which no anomalies can be present, so that there isn't a risk of the dump failing or causing other transactions to roll back with a serialization_failure. See Chapter 13 for more information about transaction isolation and concurrency control.

This option is not beneficial for a dump which is intended only for disaster recovery. It could be useful for a dump used to load a copy of the database for reporting or other read-only load sharing while the original database continues to be updated. Without it the dump may reflect a state which is not consistent with any serial execution of the transactions eventually committed. For example, if batch processing techniques are used, a batch may show as closed in the dump without all of the items which are in the batch appearing.

This option will make no difference if there are no read-write transactions active when pg_dump is started. If read-write transactions are active, the start of the dump may be delayed for an indeterminate length of time. Once running, performance with or without the switch is the same.

--snapshot=snapshotname

Use the specified synchronized snapshot when making a dump of the database (see Table 9.82 for more details).

This option is useful when needing to synchronize the dump with a logical replication slot (see Chapter 49) or with a concurrent session.

In the case of a parallel dump, the snapshot name defined by this option is used rather than taking a new snapshot.

--strict-names

Require that each schema (\-n/\--schema) and table (\-t/\--table) qualifier match at least one schema/table in the database to be dumped. Note that if none of the schema/table qualifiers find matches, pg_dump will generate an error even without --strict-names.

This option has no effect on \-N/\--exclude-schema, \-T/\--exclude-table, or \--exclude-table-data. An exclude pattern failing to match any objects is not considered an error.

--use-set-session-authorization

Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards-compatible, but depending on the history of the objects in the dump, might not restore properly. Also, a dump using SET SESSION AUTHORIZATION will certainly require superuser privileges to restore correctly, whereas ALTER OWNER requires lesser privileges.

-?
--help

Show help about pg_dump command line arguments, and exit.

The following command-line options control the database connection parameters.

-\d\ dbname
--dbname=dbname

Specifies the name of the database to connect to. This is equivalent to specifying dbname as the first non-option argument on the command line.

If this parameter contains an = sign or starts with a valid URI prefix (postgresql:// or postgres//:/), it is treated as a conninfo string. See Section 34.1 for more information.
-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_dump to prompt for a password before connecting to a database.

This option is never essential, since pg_dump will automatically prompt for a password if the server demands password authentication. However, pg_dump will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--role=rolename

Specifies a role name to be used to create the dump. This option causes pg_dump to issue a SET ROLE rolename command after connecting to the database. It is useful when the authenticated user (specified by -U) lacks privileges needed by pg_dump, but can switch to a role with the required rights. Some installations have a policy against logging in directly as a superuser, and use of this option allows dumps to be made without violating the policy.

Environment

PGDATABASE
PGHOST
PGOPTIONS
PGPORT
PGUSER

Default connection parameters.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Diagnostics

pg_dump internally executes SELECT statements. If you have problems running pg_dump, make sure you are able to select information from the database using, for example, psql. Also, any default connection settings and environment variables used by the libpq front-end library will apply.
The database activity of `pg_dump` is normally collected by the statistics collector. If this is undesirable, you can set parameter `track_counts` to false via `PGOPTIONS` or the `ALTER USER` command.

**Notes**

If your database cluster has any local additions to the `template1` database, be careful to restore the output of `pg_dump` into a truly empty database; otherwise you are likely to get errors due to duplicate definitions of the added objects. To make an empty database without any local additions, copy from `template0` not `template1`, for example:

```
CREATE DATABASE foo WITH TEMPLATE template0;
```

When a data-only dump is chosen and the option `--disable-triggers` is used, `pg_dump` emits commands to disable triggers on user tables before inserting the data, and then commands to re-enable them after the data has been inserted. If the restore is stopped in the middle, the system catalogs might be left in the wrong state.

The dump file produced by `pg_dump` does not contain the statistics used by the optimizer to make query planning decisions. Therefore, it is wise to run `ANALYZE` after restoring from a dump file to ensure optimal performance; see Section 24.1.3 and Section 24.1.6 for more information.

Because `pg_dump` is used to transfer data to newer versions of PostgreSQL, the output of `pg_dump` can be expected to load into PostgreSQL server versions newer than `pg_dump`'s version. `pg_dump` can also dump from PostgreSQL servers older than its own version. (Currently, servers back to version 8.0 are supported.) However, `pg_dump` cannot dump from PostgreSQL servers newer than its own major version; it will refuse to even try, rather than risk making an invalid dump. Also, it is not guaranteed that `pg_dump`'s output can be loaded into a server of an older major version — not even if the dump was taken from a server of that version. Loading a dump file into an older server may require manual editing of the dump file to remove syntax not understood by the older server. Use of the `--quote-all-identifiers` option is recommended in cross-version cases, as it can prevent problems arising from varying reserved-word lists in different PostgreSQL versions.

When dumping logical replication subscriptions, `pg_dump` will generate `CREATE SUBSCRIPTION` commands that use the `NOCONNECT` option, so that restoring the subscription does not make remote connections for creating a replication slot or for initial table copy. That way, the dump can be restored without requiring network access to the remote servers. It is then up to the user to reactivate the subscriptions in a suitable way. If the involved hosts have changed, the connection information might have to be changed. It might also be appropriate to truncate the target tables before initiating a new full table copy.

**Examples**

To dump a database called `mydb` into a SQL-script file:

```
$ pg_dump mydb > db.sql
```

To reload such a script into a (freshly created) database named `newdb`:

```
$ psql -d newdb -f db.sql
```

To dump a database into a custom-format archive file:

```
$ pg_dump -Fc mydb > db.dump
```

To dump a database into a directory-format archive:
$ pg_dump -Fd mydb -f dumpdir

To dump a database into a directory-format archive in parallel with 5 worker jobs:

$ pg_dump -Fd mydb -j 5 -f dumpdir

To reload an archive file into a (freshly created) database named newdb:

$ pg_restore -d newdb db.dump

To reload an archive file into the same database it was dumped from, discarding the current contents of that database:

$ pg_restore -d postgres --clean --create db.dump

To dump a single table named mytab:

$ pg_dump -t mytab mydb > db.sql

To dump all tables whose names start with emp in the detroit schema, except for the table named employee_log:

$ pg_dump -t 'detroit.emp*' -T detroit.employee_log mydb > db.sql

To dump all schemas whose names start with east or west and end in gsm, excluding any schemas whose names contain the word test:

$ pg_dump -n 'east*gsm' -n 'west*gsm' -N '*test*' mydb > db.sql

The same, using regular expression notation to consolidate the switches:

$ pg_dump -n '(east|west)*gsm' -N '*test*' mydb > db.sql

To dump all database objects except for tables whose names begin with ts_:

$ pg_dump -T 'ts_*' mydb > db.sql

To specify an upper-case or mixed-case name in -t and related switches, you need to double-quote the name; else it will be folded to lower case (see Patterns). But double quotes are special to the shell, so in turn they must be quoted. Thus, to dump a single table with a mixed-case name, you need something like

$ pg_dump -t ""MixedCaseName"" mydb > mytab.sql

See Also
pg_dumpall, pg_restore, psql
pg_dumpall

pg_dumpall — extract a PostgreSQL database cluster into a script file

Synopsis

pg_dumpall [connection-option...] [option...]

Description

pg_dumpall is a utility for writing out ("dumping") all PostgreSQL databases of a cluster into one script file. The script file contains SQL commands that can be used as input to psql to restore the databases. It does this by calling pg_dump for each database in the cluster. pg_dumpall also dumps global objects that are common to all databases, that is, database roles and tablespaces. (pg_dump does not save these objects.)

Since pg_dumpall reads tables from all databases you will most likely have to connect as a database superuser in order to produce a complete dump. Also you will need superuser privileges to execute the saved script in order to be allowed to add roles and create databases.

The SQL script will be written to the standard output. Use the [-f]file] option or shell operators to redirect it into a file.

pg_dumpall needs to connect several times to the PostgreSQL server (once per database). If you use password authentication it will ask for a password each time. It is convenient to have a ~/.pgpass file in such cases. See Section 34.15 for more information.

Options

The following command-line options control the content and format of the output.

-a
  --data-only

  Dump only the data, not the schema (data definitions).

-c
  --clean

  Include SQL commands to clean (drop) databases before recreating them. DROP commands for roles and tablespaces are added as well.

-E encoding
  --encoding=encoding

  Create the dump in the specified character set encoding. By default, the dump is created in the database encoding. (Another way to get the same result is to set the PGCLIENTENCODING environment variable to the desired dump encoding.)

-f filename
  --file=filename

  Send output to the specified file. If this is omitted, the standard output is used.

-g
  --globals-only

  Dump only global objects (roles and tablespaces), no databases.
pg_dumpall

- o
--oids
Dump object identifiers (OIDs) as part of the data for every table. Use this option if your application references the OID columns in some way (e.g., in a foreign key constraint). Otherwise, this option should not be used.

- O
--no-owner
Do not output commands to set ownership of objects to match the original database. By default, pg_dumpall issues ALTER OWNER or SET SESSION AUTHORIZATION statements to set ownership of created schema elements. These statements will fail when the script is run unless it is started by a superuser (or the same user that owns all of the objects in the script). To make a script that can be restored by any user, but will give that user ownership of all the objects, specify --no-owner.

- r
--roles-only
Dump only roles, no databases or tablespaces.

- s
--schema-only
Dump only the object definitions (schema), not data.

- S username
--superuser=username
Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used. (Usually, it's better to leave this out, and instead start the resulting script as superuser.)

- t
--tablespaces-only
Dump only tablespaces, no databases or roles.

- v
--verbose
Specifies verbose mode. This will cause pg_dumpall to output start/stop times to the dump file, and progress messages to standard error. It will also enable verbose output in pg_dump.

- V
--version
Print the pg_dumpall version and exit.

- x
--no-privileges
--no-acl
Prevent dumping of access privileges (grant/revoke commands).

--binary-upgrade
This option is for use by in-place upgrade utilities. Its use for other purposes is not recommended or supported. The behavior of the option may change in future releases without notice.
pg_dumpall

--column-inserts
--attribute-inserts

Dump data as INSERT commands with explicit column names (INSERT INTO table (column, ... ) VALUES ...). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases.

--disable-dollar-quoting

This option disables the use of dollar quoting for function bodies, and forces them to be quoted using SQL standard string syntax.

--disable-triggers

This option is relevant only when creating a data-only dump. It instructs pg_dumpall to include commands to temporarily disable triggers on the target tables while the data is reloaded. Use this if you have referential integrity checks or other triggers on the tables that you do not want to invoke during data reload.

Presently, the commands emitted for --disable-triggers must be done as superuser. So, you should also specify a superuser name with -S, or preferably be careful to start the resulting script as a superuser.

--if-exists

Use conditional commands (i.e. add an IF EXISTS clause) to drop databases and other objects. This option is not valid unless --clean is also specified.

--inserts

Dump data as INSERT commands (rather than COPY). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases. Note that the restore might fail altogether if you have rearranged column order. The --column-inserts option is safer, though even slower.

--load-via-partition-root

When dumping data for a table partition, make the COPY or INSERT statements target the root of the partitioning hierarchy that contains it, rather than the partition itself. This causes the appropriate partition to be re-determined for each row when the data is loaded. This may be useful when reloading data on a server where rows do not always fall into the same partitions as they did on the original server. That could happen, for example, if the partitioning column is of type text and the two systems have different definitions of the collation used to sort the partitioning column.

--lock-wait-timeout=timeout

Do not wait forever to acquire shared table locks at the beginning of the dump. Instead, fail if unable to lock a table within the specified timeout. The timeout may be specified in any of the formats accepted by SET statement_timeout. Allowed values vary depending on the server version you are dumping from, but an integer number of milliseconds is accepted by all versions since 7.3. This option is ignored when dumping from a pre-7.3 server.

--no-comments

Do not dump comments.

--no-publications

Do not dump publications.

--no-role-passwords

Do not dump passwords for roles. When restored, roles will have a null password, and password authentication will always fail until the password is set. Since password values aren't needed when
this option is specified, the role information is read from the catalog view pg_roles instead of pg_authid. Therefore, this option also helps if access to pg_authid is restricted by some security policy.

--no-security-labels
Do not dump security labels.

--no-subscriptions
Do not dump subscriptions.

--no-sync
By default, pg_dumpall will wait for all files to be written safely to disk. This option causes pg_dumpall to return without waiting, which is faster, but means that a subsequent operating system crash can leave the dump corrupt. Generally, this option is useful for testing but should not be used when dumping data from production installation.

--no-tablespaces
Do not output commands to create tablespaces nor select tablespaces for objects. With this option, all objects will be created in whichever tablespace is the default during restore.

--no-unlogged-table-data
Do not dump the contents of unlogged tables. This option has no effect on whether or not the table definitions (schema) are dumped; it only suppresses dumping the table data.

--quote-all-identifiers
Force quoting of all identifiers. This option is recommended when dumping a database from a server whose PostgreSQL major version is different from pg_dumpall’s, or when the output is intended to be loaded into a server of a different major version. By default, pg_dumpall quotes only identifiers that are reserved words in its own major version. This sometimes results in compatibility issues when dealing with servers of other versions that may have slightly different sets of reserved words. Using --quote-all-identifiers prevents such issues, at the price of a harder-to-read dump script.

--use-set-session-authorization
Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards compatible, but depending on the history of the objects in the dump, might not restore properly.

-?
--help
Show help about pg_dumpall command line arguments, and exit.

The following command-line options control the database connection parameters.

-d connstr
--dbname=connstr
Specifies parameters used to connect to the server, as a connection string. See Section 34.1.1 for more information.

The option is called --dbname for consistency with other client applications, but because pg_dumpall needs to connect to many databases, the database name in the connection string will be ignored. Use the -l option to specify the name of the database used for the initial connection, which will dump global objects and discover what other databases should be dumped.
pg_dumpall

- `h host`
  --host=host
  Specifies the host name of the machine on which the database server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

- `l dbname`
  --database=dbname
  Specifies the name of the database to connect to for dumping global objects and discovering what other databases should be dumped. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.

- `p port`
  --port=port
  Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

- `U username`
  --username=username
  User name to connect as.

- `w`
  --no-password
  Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

- `W`
  --password
  Force pg_dumpall to prompt for a password before connecting to a database.
  This option is never essential, since pg_dumpall will automatically prompt for a password if the server demands password authentication. However, pg_dumpall will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

  Note that the password prompt will occur again for each database to be dumped. Usually, it's better to set up a ~/.pgpass file than to rely on manual password entry.

- `--role=rolename`
  Specifies a role name to be used to create the dump. This option causes pg_dumpall to issue a SET ROLE rolename command after connecting to the database. It is useful when the authenticated user (specified by -U) lacks privileges needed by pg_dumpall, but can switch to a role with the required rights. Some installations have a policy against logging in directly as a superuser, and use of this option allows dumps to be made without violating the policy.

**Environment**

PGHOST
PGOPTIONS
PGPORT
PGUSER

Default connection parameters
This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

**Notes**

Since `pg_dumpall` calls `pg_dump` internally, some diagnostic messages will refer to `pg_dump`.

The `--clean` option can be useful even when your intention is to restore the dump script into a fresh cluster. Use of `--clean` authorizes the script to drop and re-create the built-in `postgres` and `template1` databases, ensuring that those databases will retain the same properties (for instance, locale and encoding) that they had in the source cluster. Without the option, those databases will retain their existing database-level properties, as well as any pre-existing contents.

Once restored, it is wise to run `ANALYZE` on each database so the optimizer has useful statistics. You can also run `vacuumdb -a -z` to analyze all databases.

The dump script should not be expected to run completely without errors. In particular, because the script will issue `CREATE ROLE` for every role existing in the source cluster, it is certain to get a “role already exists” error for the bootstrap superuser, unless the destination cluster was initialized with a different bootstrap superuser name. This error is harmless and should be ignored. Use of the `--clean` option is likely to produce additional harmless error messages about non-existent objects, although you can minimize those by adding `--if-exists`.

`pg_dumpall` requires all needed tablespace directories to exist before the restore; otherwise, database creation will fail for databases in non-default locations.

**Examples**

To dump all databases:

```
$ pg_dumpall > db.out
```

To reload database(s) from this file, you can use:

```
$ psql -f db.out postgres
```

It is not important to which database you connect here since the script file created by `pg_dumpall` will contain the appropriate commands to create and connect to the saved databases. An exception is that if you specified `--clean`, you must connect to the `postgres` database initially; the script will attempt to drop other databases immediately, and that will fail for the database you are connected to.

**See Also**

Check `pg_dump` for details on possible error conditions.
pg_isready

pg_isready — check the connection status of a PostgreSQL server

Synopsis

pg_isready [connection-option...] [option...]

Description

pg_isready is a utility for checking the connection status of a PostgreSQL database server. The exit status specifies the result of the connection check.

Options

- \( \text{d dbname} \)
  --dbname=dbname

  Specifies the name of the database to connect to.

  If this parameter contains an = sign or starts with a valid URI prefix (postgresql:// or postgres://), it is treated as a conninfo string. See Section 34.1.1 for more information.

- \( \text{h hostname} \)
  --host=hostname

  Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix-domain socket.

- \( \text{p port} \)
  --port=port

  Specifies the TCP port or the local Unix-domain socket file extension on which the server is listening for connections. Defaults to the value of the PGPORT environment variable or, if not set, to the port specified at compile time, usually 5432.

- \( \text{q} \)
  --quiet

  Do not display status message. This is useful when scripting.

- \( \text{t seconds} \)
  --timeout=seconds

  The maximum number of seconds to wait when attempting connection before returning that the server is not responding. Setting to 0 disables. The default is 3 seconds.

- \( \text{U username} \)
  --username=username

  Connect to the database as the user username instead of the default.

- \( \text{V} \)
  --version

  Print the pg_isready version and exit.
?-  
--help

Show help about pg_isready command line arguments, and exit.

Exit Status

pg_isready returns 0 to the shell if the server is accepting connections normally, 1 if the server is rejecting connections (for example during startup), 2 if there was no response to the connection attempt, and 3 if no attempt was made (for example due to invalid parameters).

Environment

pg_isready, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Notes

It is not necessary to supply correct user name, password, or database name values to obtain the server status; however, if incorrect values are provided, the server will log a failed connection attempt.

Examples

Standard Usage:

$ pg_isready
/tmp:5432 - accepting connections
$ echo $? 0

Running with connection parameters to a PostgreSQL cluster in startup:

$ pg_isready -h localhost -p 5433
localhost:5433 - rejecting connections
$ echo $? 1

Running with connection parameters to a non-responsive PostgreSQL cluster:

$ pg_isready -h someremotehost
someremotehost:5432 - no response
$ echo $? 2
pg_receivewal

pg_receivewal — stream write-ahead logs from a PostgreSQL server

Synopsis

pg_receivewal [option...]

Description

pg_receivewal is used to stream the write-ahead log from a running PostgreSQL cluster. The write-ahead log is streamed using the streaming replication protocol, and is written to a local directory of files. This directory can be used as the archive location for doing a restore using point-in-time recovery (see Section 25.3).

pg_receivewal streams the write-ahead log in real time as it’s being generated on the server, and does not wait for segments to complete like archive_command does. For this reason, it is not necessary to set archive_timeout when using pg_receivewal.

Unlike the WAL receiver of a PostgreSQL standby server, pg_receivewal by default flushes WAL data only when a WAL file is closed. The option --synchronous must be specified to flush WAL data in real time.

The write-ahead log is streamed over a regular PostgreSQL connection and uses the replication protocol. The connection must be made with a superuser or a user having REPLICATION permissions (see Section 21.2), and pg_hba.conf must permit the replication connection. The server must also be configured with max_wal_senders set high enough to leave at least one session available for the stream.

If the connection is lost, or if it cannot be initially established, with a non-fatal error, pg_receivewal will retry the connection indefinitely, and reestablish streaming as soon as possible. To avoid this behavior, use the -n parameter.

In the absence of fatal errors, pg_receivewal will run until terminated by the SIGINT signal (Control+C).

Options

-D directory
--directory=directory

Directory to write the output to.

This parameter is required.

-E lsn
--endpos=lsn

Automatically stop replication and exit with normal exit status 0 when receiving reaches the specified LSN.

If there is a record with LSN exactly equal to lsn, the record will be processed.

--if-not-exists

Do not error out when --create-slot is specified and a slot with the specified name already exists.
-n
--no-loop

Don't loop on connection errors. Instead, exit right away with an error.

--no-sync

This option causes `pg_receivewal` to not force WAL data to be flushed to disk. This is faster, but means that a subsequent operating system crash can leave the WAL segments corrupt. Generally, this option is useful for testing but should not be used when doing WAL archiving on a production deployment.

This option is incompatible with `--synchronous`.

-s interval
--status-interval=interval

Specifies the number of seconds between status packets sent back to the server. This allows for easier monitoring of the progress from server. A value of zero disables the periodic status updates completely, although an update will still be sent when requested by the server, to avoid timeout disconnect. The default value is 10 seconds.

-S slotname
--slot=slotname

Require `pg_receivewal` to use an existing replication slot (see Section 26.2.6). When this option is used, `pg_receivewal` will report a flush position to the server, indicating when each segment has been synchronized to disk so that the server can remove that segment if it is not otherwise needed.

When the replication client of `pg_receivewal` is configured on the server as a synchronous standby, then using a replication slot will report the flush position to the server, but only when a WAL file is closed. Therefore, that configuration will cause transactions on the primary to wait for a long time and effectively not work satisfactorily. The option `--synchronous` (see below) must be specified in addition to make this work correctly.

--synchronous

Flush the WAL data to disk immediately after it has been received. Also send a status packet back to the server immediately after flushing, regardless of `--status-interval`.

This option should be specified if the replication client of `pg_receivewal` is configured on the server as a synchronous standby, to ensure that timely feedback is sent to the server.

-v
--verbose

Enables verbose mode.

-Z level
--compress=level

Enables gzip compression of write-ahead logs, and specifies the compression level (0 through 9, 0 being no compression and 9 being best compression). The suffix `.gz` will automatically be added to all filenames.

The following command-line options control the database connection parameters.

-d connstr
--dbname=connstr

Specifies parameters used to connect to the server, as a connection string. See Section 34.1.1 for more information.
The option is called \texttt{--dbname} for consistency with other client applications, but because \texttt{pg_receivewal} doesn't connect to any particular database in the cluster, database name in the connection string will be ignored.

\texttt{-h host}  
\texttt{--host=host}  

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the \texttt{PGHOST} environment variable, if set, else a Unix domain socket connection is attempted.

\texttt{-p port}  
\texttt{--port=port}  

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the \texttt{PGPORT} environment variable, if set, or a compiled-in default.

\texttt{-U username}  
\texttt{--username=username}  

User name to connect as.

\texttt{-w}  
\texttt{--no-password}  

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a \texttt{.pgpass} file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

\texttt{-W}  
\texttt{--password}  

Force \texttt{pg_receivewal} to prompt for a password before connecting to a database.

This option is never essential, since \texttt{pg_receivewal} will automatically prompt for a password if the server demands password authentication. However, \texttt{pg_receivewal} will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing \texttt{-W} to avoid the extra connection attempt.

\texttt{pg_receivewal} can perform one of the two following actions in order to control physical replication slots:

\texttt{--create-slot}  

Create a new physical replication slot with the name specified in \texttt{--slot}, then exit.

\texttt{--drop-slot}  

Drop the replication slot with the name specified in \texttt{--slot}, then exit.

Other options are also available:

\texttt{-V}  
\texttt{--version}  

Print the \texttt{pg_receivewal} version and exit.

\texttt{-?}  
\texttt{--help}  

Show help about \texttt{pg_receivewal} command line arguments, and exit.
Exit Status

pg_receivewal will exit with status 0 when terminated by the SIGINT signal. (That is the normal way to end it. Hence it is not an error.) For fatal errors or other signals, the exit status will be nonzero.

Environment

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq (see Section 34.14).

Notes

When using pg_receivewal instead of archive_command as the main WAL backup method, it is strongly recommended to use replication slots. Otherwise, the server is free to recycle or remove write-ahead log files before they are backed up, because it does not have any information, either from archive_command or the replication slots, about how far the WAL stream has been archived. Note, however, that a replication slot will fill up the server's disk space if the receiver does not keep up with fetching the WAL data.

pg_receivewal will preserve group permissions on the received WAL files if group permissions are enabled on the source cluster.

Examples

To stream the write-ahead log from the server at mydbserver and store it in the local directory /usr/local/pgsql/archive:

```
$ pg_receivewal -h mydbserver -D /usr/local/pgsql/archive
```

See Also

pg_basebackup
pg_recvlogical

pg_recvlogical — control PostgreSQL logical decoding streams

Synopsis

pg_recvlogical [option...]

Description

pg_recvlogical controls logical decoding replication slots and streams data from such replication slots.

It creates a replication-mode connection, so it is subject to the same constraints as pg_recvewal, plus those for logical replication (see Chapter 49).

pg_recvlogical has no equivalent to the logical decoding SQL interface's peek and get modes. It sends replay confirmations for data lazily as it receives it and on clean exit. To examine pending data on a slot without consuming it, use pg_logical_slot_peek_changes.

Options

At least one of the following options must be specified to select an action:

--create-slot

Create a new logical replication slot with the name specified by --slot, using the output plugin specified by --plugin, for the database specified by --dbname.

--drop-slot

Drop the replication slot with the name specified by --slot, then exit.

--start

Begin streaming changes from the logical replication slot specified by --slot, continuing until terminated by a signal. If the server side change stream ends with a server shutdown or disconnect, retry in a loop unless --no-loop is specified.

The stream format is determined by the output plugin specified when the slot was created.

The connection must be to the same database used to create the slot.

--create-slot and --start can be specified together. --drop-slot cannot be combined with another action.

The following command-line options control the location and format of the output and other replication behavior:

-E lsn
--endpos=lsn

In --start mode, automatically stop replication and exit with normal exit status 0 when receiving reaches the specified LSN. If specified when not in --start mode, an error is raised.

If there's a record with LSN exactly equal to lsn, the record will be output.

The --endpos option is not aware of transaction boundaries and may truncate output partway through a transaction. Any partially output transaction will not be consumed and will be replayed again when the slot is next read from. Individual messages are never truncated.
-f filename
--file=filename

Write received and decoded transaction data into this file. Use - for stdout.

-F interval_seconds
--fsync-interval=interval_seconds

Specifies how often pg_recvlogical should issue fsync() calls to ensure the output file is safely flushed to disk.

The server will occasionally request the client to perform a flush and report the flush position to the server. This setting is in addition to that, to perform flushes more frequently.

Specifying an interval of 0 disables issuing fsync() calls altogether, while still reporting progress to the server. In this case, data could be lost in the event of a crash.

-I lsn
--startpos=lsn

In --start mode, start replication from the given LSN. For details on the effect of this, see the documentation in Chapter 49 and Section 53.4. Ignored in other modes.

--if-not-exists

Do not error out when --create-slot is specified and a slot with the specified name already exists.

-n
--no-loop

When the connection to the server is lost, do not retry in a loop, just exit.

-o name=value
--option=name=value

Pass the option name to the output plugin with, if specified, the option value value. Which options exist and their effects depends on the used output plugin.

-P plugin
--plugin=plugin

When creating a slot, use the specified logical decoding output plugin. See Chapter 49. This option has no effect if the slot already exists.

-s interval_seconds
--status-interval=interval_seconds

This option has the same effect as the option of the same name in pg_receivewal. See the description there.

-S slot_name
--slot=slot_name

In --start mode, use the existing logical replication slot named slot_name. In --create-slot mode, create the slot with this name. In --drop-slot mode, delete the slot with this name.

-v
--verbose

Enables verbose mode.
The following command-line options control the database connection parameters.

-\texttt{d} \textit{database}  
\texttt{--dbname=database}

The database to connect to. See the description of the actions for what this means in detail. This can be a libpq connection string; see Section 34.1.1 for more information. Defaults to user name.

-\texttt{h} \textit{hostname-or-ip}  
\texttt{--host=hostname-or-ip}

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the \texttt{PGHOST} environment variable, if set, else a Unix domain socket connection is attempted.

-\texttt{p} \textit{port}  
\texttt{--port=port}

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the \texttt{PGPORT} environment variable, if set, or a compiled-in default.

-\texttt{U} \textit{user}  
\texttt{--username=\textit{user}}

User name to connect as. Defaults to current operating system user name.

-\texttt{w}  
\texttt{--no-password}

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a \texttt{.pgpass} file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-\texttt{W}  
\texttt{--password}

Force \texttt{pg_recvlogical} to prompt for a password before connecting to a database.

This option is never essential, since \texttt{pg_recvlogical} will automatically prompt for a password if the server demands password authentication. However, \texttt{pg_recvlogical} will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing \texttt{\texttt{W}} to avoid the extra connection attempt.

The following additional options are available:

-\texttt{V}  
\texttt{--version}

Print the \texttt{pg_recvlogical} version and exit.

-\texttt{?}  
\texttt{--help}

Show help about \texttt{pg_recvlogical} command line arguments, and exit.

\section*{Environment}

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq (see Section 34.14).
Notes

pg_recvlogical will preserve group permissions on the received WAL files if group permissions are enabled on the source cluster.

Examples

See Section 49.1 for an example.

See Also

pg_receivewal
**pg_restore**

pg_restore — restore a PostgreSQL database from an archive file created by pg_dump

**Synopsis**

`pg_restore [connection-option...] [option...] [filename]`

**Description**

pg_restore is a utility for restoring a PostgreSQL database from an archive created by pg_dump in one of the non-plain-text formats. It will issue the commands necessary to reconstruct the database to the state it was in at the time it was saved. The archive files also allow pg_restore to be selective about what is restored, or even to reorder the items prior to being restored. The archive files are designed to be portable across architectures.

pg_restore can operate in two modes. If a database name is specified, pg_restore connects to that database and restores archive contents directly into the database. Otherwise, a script containing the SQL commands necessary to rebuild the database is created and written to a file or standard output. This script output is equivalent to the plain text output format of pg_dump. Some of the options controlling the output are therefore analogous to pg_dump options.

Obviously, pg_restore cannot restore information that is not present in the archive file. For instance, if the archive was made using the "dump data as INSERT commands" option, pg_restore will not be able to load the data using COPY statements.

**Options**

pg_restore accepts the following command line arguments.

*filename*

Specifies the location of the archive file (or directory, for a directory-format archive) to be restored. If not specified, the standard input is used.

-a

--data-only

Restore only the data, not the schema (data definitions). Table data, large objects, and sequence values are restored, if present in the archive.

This option is similar to, but for historical reasons not identical to, specifying --section=data.

-c

--clean

Clean (drop) database objects before recreating them. (Unless --if-exists is used, this might generate some harmless error messages, if any objects were not present in the destination database.)

-C

--create

Create the database before restoring into it. If --clean is also specified, drop and recreate the target database before connecting to it.
With `--create`, pg_restore also restores the database's comment if any, and any configuration variable settings that are specific to this database, that is, any \texttt{ALTER DATABASE ... SET ...} and \texttt{ALTER ROLE ... IN DATABASE ... SET ...} commands that mention this database. Access privileges for the database itself are also restored, unless `--no-acl` is specified.

When this option is used, the database named with `\texttt{-d}` is used only to issue the initial \texttt{DROP DATABASE} and \texttt{CREATE DATABASE} commands. All data is restored into the database name that appears in the archive.

\textbf{-d \textit{dbname}}
\textbf{--dbname=\textit{dbname}}

Connect to database \textit{dbname} and restore directly into the database.

\textbf{-e}
\textbf{--exit-on-error}

Exit if an error is encountered while sending SQL commands to the database. The default is to continue and to display a count of errors at the end of the restoration.

\textbf{-f \textit{filename}}
\textbf{--file=\textit{filename}}

Specify output file for generated script, or for the listing when used with `\texttt{-l}`. Default is the standard output.

\textbf{-F \textit{format}}
\textbf{--format=\textit{format}}

Specify format of the archive. It is not necessary to specify the format, since pg_restore will determine the format automatically. If specified, it can be one of the following:

\begin{itemize}
  \item \texttt{c} \\
  \texttt{custom}
\end{itemize}

The archive is in the custom format of pg_dump.

\begin{itemize}
  \item \texttt{d} \\
  \texttt{directory}
\end{itemize}

The archive is a directory archive.

\begin{itemize}
  \item \texttt{t} \\
  \texttt{tar}
\end{itemize}

The archive is a \texttt{tar} archive.

\textbf{-I \textit{index}}
\textbf{--index=\textit{index}}

Restore definition of named index only. Multiple indexes may be specified with multiple `-I` switches.

\textbf{-j \textit{number-of-jobs}}
\textbf{--jobs=\textit{number-of-jobs}}

Run the most time-consuming parts of pg_restore — those which load data, create indexes, or create constraints — using multiple concurrent jobs. This option can dramatically reduce the time to restore a large database to a server running on a multiprocessor machine.

Each job is one process or one thread, depending on the operating system, and uses a separate connection to the server.
The optimal value for this option depends on the hardware setup of the server, of the client, and of the network. Factors include the number of CPU cores and the disk setup. A good place to start is the number of CPU cores on the server, but values larger than that can also lead to faster restore times in many cases. Of course, values that are too high will lead to decreased performance because of thrashing.

Only the custom and directory archive formats are supported with this option. The input must be a regular file or directory (not, for example, a pipe). This option is ignored when emitting a script rather than connecting directly to a database server. Also, multiple jobs cannot be used together with the option `--single-transaction`.

```
-l
--list
```

List the table of contents of the archive. The output of this operation can be used as input to the `--list` option. Note that if filtering switches such as `--n` or `--t` are used with `--l`, they will restrict the items listed.

```
-L list-file
--use-list=list-file
```

Restore only those archive elements that are listed in `list-file`, and restore them in the order they appear in the file. Note that if filtering switches such as `--n` or `--t` are used with `--L`, they will further restrict the items restored.

`list-file` is normally created by editing the output of a previous `--l` operation. Lines can be moved or removed, and can also be commented out by placing a semicolon (`;`) at the start of the line. See below for examples.

```
-n schema
--schema=schema
```

Restore only objects that are in the named schema. Multiple schemas may be specified with multiple `--n` switches. This can be combined with the `--t` option to restore just a specific table.

```
-N schema
--exclude-schema=schema
```

Do not restore objects that are in the named schema. Multiple schemas to be excluded may be specified with multiple `--N` switches.

When both `--n` and `--N` are given for the same schema name, the `--N` switch wins and the schema is excluded.

```
-O
--no-owner
```

Do not output commands to set ownership of objects to match the original database. By default, `pg_restore` issues `ALTER OWNER` or `SET SESSION AUTHORIZATION` statements to set ownership of created schema elements. These statements will fail unless the initial connection to the database is made by a superuser (or the same user that owns all of the objects in the script). With `--O`, any user name can be used for the initial connection, and this user will own all the created objects.

```
-P function-name(argtype [, ...])
--function=function-name(argtype [, ...])
```

Restore the named function only. Be careful to spell the function name and arguments exactly as they appear in the dump file's table of contents. Multiple functions may be specified with multiple `--P` switches.
-R
--no-reconnect

This option is obsolete but still accepted for backwards compatibility.

-s
--schema-only

Restore only the schema (data definitions), not data, to the extent that schema entries are present in the archive.

This option is the inverse of --data-only. It is similar to, but for historical reasons not identical to, specifying --section=pre-data --section=post-data.

(Do not confuse this with the --schema option, which uses the word “schema” in a different meaning.)

-S username
--superuser=username

Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used.

-t table
--table=table

Restore definition and/or data of only the named table. For this purpose, “table” includes views, materialized views, sequences, and foreign tables. Multiple tables can be selected by writing multiple -t switches. This option can be combined with the -n option to specify table(s) in a particular schema.

**Note**

When -t is specified, pg_restore makes no attempt to restore any other database objects that the selected table(s) might depend upon. Therefore, there is no guarantee that a specific-table restore into a clean database will succeed.

**Note**

This flag does not behave identically to the -t flag of pg_dump. There is not currently any provision for wild-card matching in pg_restore, nor can you include a schema name within its -t. And, while pg_dump’s -t flag will also dump subsidiary objects (such as indexes) of the selected table(s), pg_restore’s -t flag does not include such subsidiary objects.

**Note**

In versions prior to PostgreSQL 9.6, this flag matched only tables, not any other type of relation.

-T trigger
--trigger=trigger

Restore named trigger only. Multiple triggers may be specified with multiple -T switches.
-v
--verbose

Specifies verbose mode.

-V
--version

Print the pg_restore version and exit.

-x
--no-privileges
--no-acl

Prevent restoration of access privileges (grant/revoke commands).

-1
--single-transaction

Execute the restore as a single transaction (that is, wrap the emitted commands in BEGIN/COMMIT). This ensures that either all the commands complete successfully, or no changes are applied. This option implies --exit-on-error.

--disable-triggers

This option is relevant only when performing a data-only restore. It instructs pg_restore to execute commands to temporarily disable triggers on the target tables while the data is reloaded. Use this if you have referential integrity checks or other triggers on the tables that you do not want to invoke during data reload.

Presently, the commands emitted for --disable-triggers must be done as superuser. So you should also specify a superuser name with -S or, preferably, run pg_restore as a PostgreSQL superuser.

--enable-row-security

This option is relevant only when restoring the contents of a table which has row security. By default, pg_restore will set row_security to off, to ensure that all data is restored in to the table. If the user does not have sufficient privileges to bypass row security, then an error is thrown. This parameter instructs pg_restore to set row_security to on instead, allowing the user to attempt to restore the contents of the table with row security enabled. This might still fail if the user does not have the right to insert the rows from the dump into the table.

Note that this option currently also requires the dump be in INSERT format, as COPY FROM does not support row security.

--if-exists

Use conditional commands (i.e. add an IF EXISTS clause) to drop database objects. This option is not valid unless --clean is also specified.

--no-comments

Do not output commands to restore comments, even if the archive contains them.

--no-data-for-failed-tables

By default, table data is restored even if the creation command for the table failed (e.g., because it already exists). With this option, data for such a table is skipped. This behavior is useful if the target database already contains the desired table contents. For example, auxiliary tables for
PostgreSQL extensions such as PostGIS might already be loaded in the target database; specifying this option prevents duplicate or obsolete data from being loaded into them.

This option is effective only when restoring directly into a database, not when producing SQL script output.

--no-publications
Do not output commands to restore publications, even if the archive contains them.

--no-security-labels
Do not output commands to restore security labels, even if the archive contains them.

--no-subscriptions
Do not output commands to restore subscriptions, even if the archive contains them.

--no-tablespaces
Do not output commands to select tablespaces. With this option, all objects will be created in whichever tablespace is the default during restore.

--section=sectionname
Only restore the named section. The section name can be pre-data, data, or post-data. This option can be specified more than once to select multiple sections. The default is to restore all sections.

The data section contains actual table data as well as large-object definitions. Post-data items consist of definitions of indexes, triggers, rules and constraints other than validated check constraints. Pre-data items consist of all other data definition items.

--strict-names
Require that each schema (-n|--schema) and table (-t|--table) qualifier match at least one schema/table in the backup file.

--use-set-session-authorization
Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards-compatible, but depending on the history of the objects in the dump, might not restore properly.

-?
--help

Show help about pg_restore command line arguments, and exit.

pg_restore also accepts the following command line arguments for connection parameters:

- h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

- p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.
pg_restore

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_restore to prompt for a password before connecting to a database.

This option is never essential, since pg_restore will automatically prompt for a password if the server demands password authentication. However, pg_restore will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--role=rolemame

Specifies a role name to be used to perform the restore. This option causes pg_restore to issue a SET ROLE rolemame command after connecting to the database. It is useful when the authenticated user (specified by -U) lacks privileges needed by pg_restore, but can switch to a role with the required rights. Some installations have a policy against logging in directly as a superuser, and use of this option allows restores to be performed without violating the policy.

Environment

PGHOST
PGOPTIONS
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14). However, it does not read PGDATABASE when a database name is not supplied.

Diagnostics

When a direct database connection is specified using the -d option, pg_restore internally executes SQL statements. If you have problems running pg_restore, make sure you are able to select information from the database using, for example, psql. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Notes

If your installation has any local additions to the template1 database, be careful to load the output of pg_restore into a truly empty database; otherwise you are likely to get errors due to duplicate definitions of the added objects. To make an empty database without any local additions, copy from template0 not template1, for example:

CREATE DATABASE foo WITH TEMPLATE template0;
The limitations of `pg_restore` are detailed below.

- When restoring data to a pre-existing table and the option `--disable-triggers` is used, `pg_restore` emits commands to disable triggers on user tables before inserting the data, then emits commands to re-enable them after the data has been inserted. If the restore is stopped in the middle, the system catalogs might be left in the wrong state.

- `pg_restore` cannot restore large objects selectively; for instance, only those for a specific table. If an archive contains large objects, then all large objects will be restored, or none of them if they are excluded via `--largeobject`, `--t`, or other options.

See also the `pg_dump` documentation for details on limitations of `pg_dump`.

Once restored, it is wise to run `ANALYZE` on each restored table so the optimizer has useful statistics; see Section 24.1.3 and Section 24.1.6 for more information.

### Examples

Assume we have dumped a database called `mydb` into a custom-format dump file:

```bash
$ pg_dump -Fc mydb > db.dump
```

To drop the database and recreate it from the dump:

```bash
$ dropdb mydb
$ pg_restore -C -d postgres db.dump
```

The database named in the `-d` switch can be any database existing in the cluster; `pg_restore` only uses it to issue the `CREATE DATABASE` command for `mydb`. With `-C`, data is always restored into the database name that appears in the dump file.

To reload the dump into a new database called `newdb`:

```bash
$ createdb -T template0 newdb
$ pg_restore -d newdb db.dump
```

Notice we don’t use `-C`, and instead connect directly to the database to be restored into. Also note that we clone the new database from `template0` not `template1`, to ensure it is initially empty.

To reorder database items, it is first necessary to dump the table of contents of the archive:

```bash
$ pg_restore -l db.dump > db.list
```

The listing file consists of a header and one line for each item, e,g.:
Semicolons start a comment, and the numbers at the start of lines refer to the internal archive ID assigned to each item.

Lines in the file can be commented out, deleted, and reordered. For example:

```sql
10; 145433 TABLE map_resolutions postgres
;2; 145344 TABLE species postgres
;4; 145359 TABLE nt_header postgres
6; 145402 TABLE species_records postgres
;8; 145416 TABLE ss_old postgres
```

could be used as input to `pg_restore` and would only restore items 10 and 6, in that order:

```
$ pg_restore -L db.list db.dump
```

**See Also**

`pg_dump`, `pg_dumpall`, `psql`
psql

psql — PostgreSQL interactive terminal

Synopsis

psql [option...] [dbname [username]]

Description

psql is a terminal-based front-end to PostgreSQL. It enables you to type in queries interactively, issue them to PostgreSQL, and see the query results. Alternatively, input can be from a file or from command line arguments. In addition, psql provides a number of meta-commands and various shell-like features to facilitate writing scripts and automating a wide variety of tasks.

Options

- -a
  --echo-all

  Print all nonempty input lines to standard output as they are read. (This does not apply to lines read interactively.) This is equivalent to setting the variable ECHO to all.

- -A
  --no-align

  Switches to unaligned output mode. (The default output mode is otherwise aligned.) This is equivalent to \pset format unaligned.

- -b
  --echo-errors

  Print failed SQL commands to standard error output. This is equivalent to setting the variable ECHO to errors.

- -c command
  --command=command

  Specifies that psql is to execute the given command string, command. This option can be repeated and combined in any order with the -f option. When either -c or -f is specified, psql does not read commands from standard input; instead it terminates after processing all the -c and -f options in sequence.

  command must be either a command string that is completely parsable by the server (i.e., it contains no psql-specific features), or a single backslash command. Thus you cannot mix SQL and psql meta-commands within a -c option. To achieve that, you could use repeated -c options or pipe the string into psql, for example:

  psql -c '\x' -c 'SELECT * FROM foo;'  

  or

  echo '\x \ SELECT * FROM foo;' | psql  

  (\ is the separator meta-command.)
Each SQL command string passed to \texttt{-c} is sent to the server as a single request. Because of this, the server executes it as a single transaction even if the string contains multiple SQL commands, unless there are explicit \texttt{BEGIN/COMMIT} commands included in the string to divide it into multiple transactions. (See Section 53.2.2.1 for more details about how the server handles multi-query strings.) Also, \textit{psql} only prints the result of the last SQL command in the string. This is different from the behavior when the same string is read from a file or fed to \textit{psql}'s standard input, because then \textit{psql} sends each SQL command separately.

Because of this behavior, putting more than one SQL command in a single \texttt{-c} string often has unexpected results. It's better to use repeated \texttt{-c} commands or feed multiple commands to \textit{psql}'s standard input, either using echo as illustrated above, or via a shell here-document, for example:

\begin{verbatim}
psql <<EOF
  \x
  SELECT * FROM foo;
EOF
-d dbname
--dbname=dbname
\end{verbatim}

Specifies the name of the database to connect to. This is equivalent to specifying \textit{dbname} as the first non-option argument on the command line.

If this parameter contains an = sign or starts with a valid URI prefix (postgres:// or postgres://), it is treated as a \textit{conninfo} string. See Section 34.1.1 for more information.

-\texttt{e}
--\texttt{echo-queries}

Copy all SQL commands sent to the server to standard output as well. This is equivalent to setting the variable \texttt{ECHO} to \texttt{queries}.

-\texttt{E}
--\texttt{echo-hidden}

Echo the actual queries generated by \texttt{\textbackslash d} and other backslash commands. You can use this to study \textit{psql}'s internal operations. This is equivalent to setting the variable \texttt{ECHO_HIDDEN} to \texttt{on}.

-\texttt{f} \textit{filename}
--\texttt{file=}\textit{filename}

Read commands from the file \textit{filename}, rather than standard input. This option can be repeated and combined in any order with the \texttt{-c} option. When either \texttt{-c} or \texttt{-f} is specified, \textit{psql} does not read commands from standard input; instead it terminates after processing all the \texttt{-c} and \texttt{-f} options in sequence. Except for that, this option is largely equivalent to the meta-command \texttt{\textbackslash i}.

If \textit{filename} is \texttt{-} (hyphen), then standard input is read until an EOF indication or \texttt{\textbackslash q} meta-command. This can be used to intersperse interactive input with input from files. Note however that Readline is not used in this case (much as if \texttt{-n} had been specified).

Using this option is subtly different from writing \textit{psql} \textless \textit{filename}. In general, both will do what you expect, but using \texttt{-f} enables some nice features such as error messages with line numbers. There is also a slight chance that using this option will reduce the start-up overhead.

On the other hand, the variant using the shell's input redirection is (in theory) guaranteed to yield exactly the same output you would have received had you entered everything by hand.

-\texttt{F} \textit{separator}
--\texttt{field-separator=}\textit{separator}

Use \textit{separator} as the field separator for unaligned output. This is equivalent to \texttt{\textbackslash pset fieldsep} or \texttt{\textbackslash f}.  

1899
-h hostname
--host=hostname

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix-domain socket.

-H
--html

Turn on HTML tabular output. This is equivalent to \pset format html or the \H command.

-l
--list

List all available databases, then exit. Other non-connection options are ignored. This is similar to the meta-command \list.

When this option is used, psql will connect to the database postgres, unless a different database is named on the command line (option -d or non-option argument, possibly via a service entry, but not via an environment variable).

-L filename
--log-file=filename

Write all query output into file filename, in addition to the normal output destination.

-n
--no-readline

Do not use Readline for line editing and do not use the command history. This can be useful to turn off tab expansion when cutting and pasting.

-o filename
--output=filename

Put all query output into file filename. This is equivalent to the command \o.

-p port
--port=port

Specifies the TCP port or the local Unix-domain socket file extension on which the server is listening for connections. Defaults to the value of the PGPORT environment variable or, if not set, to the port specified at compile time, usually 5432.

-P assignment
--pset=assignment

Specifies printing options, in the style of \pset. Note that here you have to separate name and value with an equal sign instead of a space. For example, to set the output format to LaTeX, you could write -P format=latex.

-q
--quiet

Specifies that psql should do its work quietly. By default, it prints welcome messages and various informational output. If this option is used, none of this happens. This is useful with the -c option. This is equivalent to setting the variable QUIET to on.

-R separator
--record-separator=separator

Use separator as the record separator for unaligned output. This is equivalent to \pset recordsep.
-s
--single-step

Run in single-step mode. That means the user is prompted before each command is sent to the server, with the option to cancel execution as well. Use this to debug scripts.

-S
--single-line

Runs in single-line mode where a newline terminates an SQL command, as a semicolon does.

```
Note
This mode is provided for those who insist on it, but you are not necessarily encouraged to use it. In particular, if you mix SQL and meta-commands on a line the order of execution might not always be clear to the inexperienced user.
```

-t
--tuples-only

Turn off printing of column names and result row count footers, etc. This is equivalent to \t or \pset tuples_only.

-T table_options
--table-attr=table_options

Specifies options to be placed within the HTML table tag. See \pset tableattr for details.

-U username
--username=username

Connect to the database as the user username instead of the default. (You must have permission to do so, of course.)

-v assignment
--set=assignment
--variable=assignment

Perform a variable assignment, like the \set meta-command. Note that you must separate name and value, if any, by an equal sign on the command line. To unset a variable, leave off the equal sign. To set a variable with an empty value, use the equal sign but leave off the value. These assignments are done during command line processing, so variables that reflect connection state will get overwritten later.

-V
--version

Print the psql version and exit.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

Note that this option will remain set for the entire session, and so it affects uses of the meta-command \connect as well as the initial connection attempt.
psql

-W
--password

Force psql to prompt for a password before connecting to a database.

This option is never essential, since psql will automatically prompt for a password if the server demands password authentication. However, psql will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

Note that this option will remain set for the entire session, and so it affects uses of the meta-command \connect as well as the initial connection attempt.

-x
--expanded

Turn on the expanded table formatting mode. This is equivalent to \x or \pset expanded.

-X,
--no-psqlrc

Do not read the start-up file (neither the system-wide psqlrc file nor the user's ~/.psqlrc file).

-z
--field-separator-zero

Set the field separator for unaligned output to a zero byte. This is equivalent to \pset fieldsep_zero.

-0
--record-separator-zero

Set the record separator for unaligned output to a zero byte. This is useful for interfacing, for example, with xargs -0. This is equivalent to \pset recordsep_zero.

-1
--single-transaction

This option can only be used in combination with one or more -c and/or -f options. It causes psql to issue a BEGIN command before the first such option and a COMMIT command after the last one, thereby wrapping all the commands into a single transaction. This ensures that either all the commands complete successfully, or no changes are applied.

If the commands themselves contain BEGIN, COMMIT, or ROLLBACK, this option will not have the desired effects. Also, if an individual command cannot be executed inside a transaction block, specifying this option will cause the whole transaction to fail.

-?
--help[=topic]

Show help about psql and exit. The optional topic parameter (defaulting to options) selects which part of psql is explained: commands describes psql's backslash commands; options describes the command-line options that can be passed to psql; and variables shows help about psql configuration variables.

**Exit Status**

psql returns 0 to the shell if it finished normally, 1 if a fatal error of its own occurs (e.g. out of memory, file not found), 2 if the connection to the server went bad and the session was not interactive, and 3 if an error occurred in a script and the variable ON_ERROR_STOP was set.
Usage

Connecting to a Database

psql is a regular PostgreSQL client application. In order to connect to a database you need to know the name of your target database, the host name and port number of the server, and what user name you want to connect as. psql can be told about those parameters via command line options, namely -d, -h, -p, and -U respectively. If an argument is found that does not belong to any option it will be interpreted as the database name (or the user name, if the database name is already given). Not all of these options are required; there are useful defaults. If you omit the host name, psql will connect via a Unix-domain socket to a server on the local host, or via TCP/IP to localhost on machines that don't have Unix-domain sockets. The default port number is determined at compile time. Since the database server uses the same default, you will not have to specify the port in most cases. The default user name is your operating-system user name, as is the default database name. Note that you cannot just connect to any database under any user name. Your database administrator should have informed you about your access rights.

When the defaults aren't quite right, you can save yourself some typing by setting the environment variables PGDATABASE, PGHOST, PGPORT and/or PGUSER to appropriate values. (For additional environment variables, see Section 34.14.) It is also convenient to have a ~/.pgpass file to avoid regularly having to type in passwords. See Section 34.15 for more information.

An alternative way to specify connection parameters is in a conninfo string or a URI, which is used instead of a database name. This mechanism give you very wide control over the connection. For example:

$ psql "service=myservice sslmode=require"
$ psql postgresql://dbmaster:5433/mydb?sslmode=require

This way you can also use LDAP for connection parameter lookup as described in Section 34.17. See Section 34.1.2 for more information on all the available connection options.

If the connection could not be made for any reason (e.g., insufficient privileges, server is not running on the targeted host, etc.), psql will return an error and terminate.

If both standard input and standard output are a terminal, then psql sets the client encoding to “auto”, which will detect the appropriate client encoding from the locale settings (LC_CTYPE environment variable on Unix systems). If this doesn't work out as expected, the client encoding can be overridden using the environment variable PGCLIENTENCODING.

Entering SQL Commands

In normal operation, psql provides a prompt with the name of the database to which psql is currently connected, followed by the string =>. For example:

$ psql testdb
psql (11.1)
Type "help" for help.

testdb=>

At the prompt, the user can type in SQL commands. Ordinarily, input lines are sent to the server when a command-terminating semicolon is reached. An end of line does not terminate a command. Thus commands can be spread over several lines for clarity. If the command was sent and executed without error, the results of the command are displayed on the screen.

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin your session by removing publicly-writable schemas from search_path. One can add op--
tions=-csearch_path= to the connection string or issue SELECT pg_catalog.set_config('search_path', '', false) before other SQL commands. This consideration is not specific to psql; it applies to every interface for executing arbitrary SQL commands.

Whenever a command is executed, psql also polls for asynchronous notification events generated by LISTEN and NOTIFY.

While C-style block comments are passed to the server for processing and removal, SQL-standard comments are removed by psql.

**Meta-Commands**

Anything you enter in psql that begins with an unquoted backslash is a psql meta-command that is processed by psql itself. These commands make psql more useful for administration or scripting. Meta-commands are often called slash or backslash commands.

The format of a psql command is the backslash, followed immediately by a command verb, then any arguments. The arguments are separated from the command verb and each other by any number of whitespace characters.

To include whitespace in an argument you can quote it with single quotes. To include a single quote in an argument, write two single quotes within single-quoted text. Anything contained in single quotes is furthermore subject to C-like substitutions for \n (new line), \t (tab), \b (backspace), \r (carriage return), \f (form feed), \digits (octal), and \x\digits (hexadecimal). A backslash preceding any other character within single-quoted text quotes that single character, whatever it is.

If an unquoted colon (:) followed by a psql variable name appears within an argument, it is replaced by the variable's value, as described in SQL Interpolation. The forms :"variable_name" and :"variable_name" described there work as well. The :{"variable_name} syntax allows testing whether a variable is defined. It is substituted by TRUE or FALSE. Escaping the colon with a backslash protects it from substitution.

Within an argument, text that is enclosed in backquotes (`) is taken as a command line that is passed to the shell. The output of the command (with any trailing newline removed) replaces the backquoted text. Within the text enclosed in backquotes, no special quoting or other processing occurs, except that appearances of :"variable_name" where variable_name is a psql variable name are replaced by the variable's value. Also, appearances of :"variable_name" are replaced by the variable's value suitably quoted to become a single shell command argument. (The latter form is almost always preferable, unless you are very sure of what is in the variable.) Because carriage return and line feed characters cannot be safely quoted on all platforms, the :"variable_name" form prints an error message and does not substitute the variable value when such characters appear in the value.

Some commands take an SQL identifier (such as a table name) as argument. These arguments follow the syntax rules of SQL: Unquoted letters are forced to lowercase, while double quotes (") protect letters from case conversion and allow incorporation of whitespace into the identifier. Within double quotes, paired double quotes reduce to a single double quote in the resulting name. For example, FOO"BAR"BAZ is interpreted as fooBARbaz, and "A weird" name becomes A weird name.

Parsing for arguments stops at the end of the line, or when another unquoted backslash is found. An unquoted backslash is taken as the beginning of a new meta-command. The special sequence \ \ (two backslashes) marks the end of arguments and continues parsing SQL commands, if any. That way SQL and psql commands can be freely mixed on a line. But in any case, the arguments of a meta-command cannot continue beyond the end of the line.

Many of the meta-commands act on the current query buffer. This is simply a buffer holding whatever SQL command text has been typed but not yet sent to the server for execution. This will include previous input lines as well as any text appearing before the meta-command on the same line.

The following meta-commands are defined:
If the current table output format is unaligned, it is switched to aligned. If it is not unaligned, it is set to unaligned. This command is kept for backwards compatibility. See \pset for a more general solution.

\c or \connect [ -reuse-previous=on|off ] [ dbname [ username ] [ host ] [ port ] | conninfo ]

Establishes a new connection to a PostgreSQL server. The connection parameters to use can be specified either using a positional syntax, or using conninfo connection strings as detailed in Section 34.1.1.

Where the command omits database name, user, host, or port, the new connection can reuse values from the previous connection. By default, values from the previous connection are reused except when processing a conninfo string. Passing a first argument of -reuse-previous=on or -reuse-previous=off overrides that default. When the command neither specifies nor reuses a particular parameter, the libpq default is used. Specifying any of dbname, username, host or port as - is equivalent to omitting that parameter.

If the new connection is successfully made, the previous connection is closed. If the connection attempt failed (wrong user name, access denied, etc.), the previous connection will only be kept if psql is in interactive mode. When executing a non-interactive script, processing will immediately stop with an error. This distinction was chosen as a user convenience against typos on the one hand, and a safety mechanism that scripts are not accidentally acting on the wrong database on the other hand.

Examples:

=> \c mydb myuser host.dom 6432
=> \c service=foo
=> \c "host=localhost port=5432 dbname=mydb connect_timeout=10 sslmode=disable"
=> \c postgresql://tom@localhost/mydb?application_name=myapp

\C [ title ]

Sets the title of any tables being printed as the result of a query or unset any such title. This command is equivalent to \pset title title. (The name of this command derives from “caption”, as it was previously only used to set the caption in an HTML table.)

\cd [ directory ]

Changes the current working directory to directory. Without argument, changes to the current user’s home directory.

Tip
To print your current working directory, use \! pwd.

\conninfo

Outputs information about the current database connection.

\copy { table [ ( column_list ) ] | ( query ) } { from | to } { 'filename' | program 'command' | stdin | stdout | pstdin | pstdout } [ [ with ] ( option [ , ... ] ) ]

Performs a frontend (client) copy. This is an operation that runs an SQL COPY command, but instead of the server reading or writing the specified file, psql reads or writes the file and routes the
data between the server and the local file system. This means that file accessibility and privileges are those of the local user, not the server, and no SQL superuser privileges are required.

When program is specified, command is executed by psql and the data passed from or to command is routed between the server and the client. Again, the execution privileges are those of the local user, not the server, and no SQL superuser privileges are required.

For \copy ... from stdin, data rows are read from the same source that issued the command, continuing until \ is read or the stream reaches EOF. This option is useful for populating tables in-line within a SQL script file. For \copy ... to stdout, output is sent to the same place as psql command output, and the COPY count command status is not printed (since it might be confused with a data row). To read/write psql's standard input or output regardless of the current command source or \o option, write from pstdin or to pstdout.

The syntax of this command is similar to that of the SQL COPY command. All options other than the data source/destination are as specified for COPY. Because of this, special parsing rules apply to the \copy meta-command. Unlike most other meta-commands, the entire remainder of the line is always taken to be the arguments of \copy, and neither variable interpolation nor backquote expansion are performed in the arguments.

Tip

This operation is not as efficient as the SQL COPY command because all data must pass through the client/server connection. For large amounts of data the SQL command might be preferable.

\copyright

Shows the copyright and distribution terms of PostgreSQL.


Executes the current query buffer (like \g) and shows the results in a crosstab grid. The query must return at least three columns. The output column identified by colV becomes a vertical header and the output column identified by colH becomes a horizontal header. colD identifies the output column to display within the grid. sortcolH identifies an optional sort column for the horizontal header.

Each column specification can be a column number (starting at 1) or a column name. The usual SQL case folding and quoting rules apply to column names. If omitted, colV is taken as column 1 and colH as column 2. colH must differ from colV. If colD is not specified, then there must be exactly three columns in the query result, and the column that is neither colV nor colH is taken to be colD.

The vertical header, displayed as the leftmost column, contains the values found in column colV, in the same order as in the query results, but with duplicates removed.

The horizontal header, displayed as the first row, contains the values found in column colH, with duplicates removed. By default, these appear in the same order as in the query results. But if the optional sortcolH argument is given, it identifies a column whose values must be integer numbers, and the values from colH will appear in the horizontal header sorted according to the corresponding sortcolH values.

Inside the crosstab grid, for each distinct value x of colH and each distinct value y of colV, the cell located at the intersection (x, y) contains the value of the colD column in the query result row for which the value of colH is x and the value of colV is y. If there is no such row, the cell is empty. If there are multiple such rows, an error is reported.
\d[S+] [ pattern ]

For each relation (table, view, materialized view, index, sequence, or foreign table) or composite type matching the pattern, show all columns, their types, the tablespace (if not the default) and any special attributes such as NOT NULL or defaults. Associated indexes, constraints, rules, and triggers are also shown. For foreign tables, the associated foreign server is shown as well. ("Matching the pattern" is defined in Patterns below.)

For some types of relation, \d shows additional information for each column: column values for sequences, indexed expressions for indexes, and foreign data wrapper options for foreign tables.

The command form \d+ is identical, except that more information is displayed: any comments associated with the columns of the table are shown, as is the presence of OIDs in the table, the view definition if the relation is a view, a non-default replica identity setting.

By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects.

Note

If \d is used without a pattern argument, it is equivalent to \dtvmsE which will show a list of all visible tables, views, materialized views, sequences and foreign tables. This is purely a convenience measure.

\da[S] [ pattern ]

Lists aggregate functions, together with their return type and the data types they operate on. If pattern is specified, only aggregates whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects.

\dA[+] [ pattern ]

Lists access methods. If pattern is specified, only access methods whose names match the pattern are shown. If + is appended to the command name, each access method is listed with its associated handler function and description.

\db[+] [ pattern ]

Lists tablespaces. If pattern is specified, only tablespaces whose names match the pattern are shown. If + is appended to the command name, each tablespace is listed with its associated options, on-disk size, permissions and description.

\dc[S+] [ pattern ]

Lists conversions between character-set encodings. If pattern is specified, only conversions whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects. If + is appended to the command name, each object is listed with its associated description.

\dC[+] [ pattern ]

Lists type casts. If pattern is specified, only casts whose source or target types match the pattern are listed. If + is appended to the command name, each object is listed with its associated description.

\dd[S] [ pattern ]

Shows the descriptions of objects of type constraint, operator class, operator family, rule, and trigger. All other comments may be viewed by the respective backslash commands for those object types.
\dd displays descriptions for objects matching the pattern, or of visible objects of the appropriate type if no argument is given. But in either case, only objects that have a description are listed. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects.

Descriptions for objects can be created with the COMMENT SQL command.

\dD[S+] [ pattern ]

Lists domains. If pattern is specified, only domains whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects. If + is appended to the command name, each object is listed with its associated permissions and description.

\ddp [ pattern ]

Lists default access privilege settings. An entry is shown for each role (and schema, if applicable) for which the default privilege settings have been changed from the built-in defaults. If pattern is specified, only entries whose role name or schema name matches the pattern are listed.

The ALTER DEFAULT PRIVILEGES command is used to set default access privileges. The meaning of the privilege display is explained under GRANT.

\dE[S+] [ pattern ]
\di[S+] [ pattern ]
\dm[S+] [ pattern ]
\ds[S+] [ pattern ]
\dt[S+] [ pattern ]
\dv[S+] [ pattern ]

In this group of commands, the letters E, i, m, s, t, and v stand for foreign table, index, materialized view, sequence, table, and view, respectively. You can specify any or all of these letters, in any order, to obtain a listing of objects of these types. For example, \dit lists indexes and tables. If + is appended to the command name, each object is listed with its physical size on disk and its associated description, if any. If pattern is specified, only objects whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects.

\des[+] [ pattern ]

Lists foreign servers (mnemonic: “external servers”). If pattern is specified, only those servers whose name matches the pattern are listed. If the form \des+ is used, a full description of each server is shown, including the server's ACL, type, version, options, and description.

\det[+] [ pattern ]

Lists foreign tables (mnemonic: “external tables”). If pattern is specified, only entries whose table name or schema name matches the pattern are listed. If the form \det+ is used, generic options and the foreign table description are also displayed.

\deu[+] [ pattern ]

Lists user mappings (mnemonic: “external users”). If pattern is specified, only those mappings whose user names match the pattern are listed. If the form \deu+ is used, additional information about each mapping is shown.

\deu+ might also display the user name and password of the remote user, so care should be taken not to disclose them.
\dew[+][ pattern ]

Lists foreign-data wrappers (mnemonic: “external wrappers”). If \textit{pattern} is specified, only those foreign-data wrappers whose name matches the pattern are listed. If the form \dew+ is used, the ACL, options, and description of the foreign-data wrapper are also shown.

\df[anptwS]+ [ pattern ]

Lists functions, together with their result data types, argument data types, and function types, which are classified as “agg” (aggregate), “normal”, “procedure”, “trigger”, or “window”. To display only functions of specific type(s), add the corresponding letters \texttt{a}, \texttt{n}, \texttt{p}, \texttt{t}, or \texttt{w} to the command. If \textit{pattern} is specified, only functions whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If the form \df+ is used, additional information about each function is shown, including volatility, parallel safety, owner, security classification, access privileges, language, source code and description.

\begin{quote}
\textbf{Tip}
To look up functions taking arguments or returning values of a specific data type, use your pager’s search capability to scroll through the \texttt{\df} output.
\end{quote}

\dF[+][ pattern ]

Lists text search configurations. If \textit{pattern} is specified, only configurations whose names match the pattern are shown. If the form \dF+ is used, a full description of each configuration is shown, including the underlying text search parser and the dictionary list for each parser token type.

\dFd[+][ pattern ]

Lists text search dictionaries. If \textit{pattern} is specified, only dictionaries whose names match the pattern are shown. If the form \dFd+ is used, additional information is shown about each selected dictionary, including the underlying text search template and the option values.

\dFp[+][ pattern ]

Lists text search parsers. If \textit{pattern} is specified, only parsers whose names match the pattern are shown. If the form \dFp+ is used, a full description of each parser is shown, including the underlying functions and the list of recognized token types.

\dFt[+][ pattern ]

Lists text search templates. If \textit{pattern} is specified, only templates whose names match the pattern are shown. If the form \dFt+ is used, additional information is shown about each template, including the underlying function names.

\dg[S]+ [ pattern ]

Lists database roles. (Since the concepts of “users” and “groups” have been unified into “roles”, this command is now equivalent to \texttt{\du}.) By default, only user-created roles are shown; supply the \texttt{S} modifier to include system roles. If \textit{pattern} is specified, only those roles whose names match the pattern are listed. If the form \dg+ is used, additional information is shown about each role; currently this adds the comment for each role.

\dl

This is an alias for \texttt{\lo_list}, which shows a list of large objects.
\dL[\+][ ]{ pattern }  
Lists procedural languages. If \texttt{pattern} is specified, only languages whose names match the pattern are listed. By default, only user-created languages are shown; supply the \texttt{S} modifier to include system objects. If \texttt{+} is appended to the command name, each language is listed with its call handler, validator, access privileges, and whether it is a system object.

\dn[\+][ ]{ pattern }  
Lists schemas (namespaces). If \texttt{pattern} is specified, only schemas whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If \texttt{+} is appended to the command name, each object is listed with its associated permissions and description, if any.

\do[\+][ ]{ pattern }  
Lists operators with their operand and result types. If \texttt{pattern} is specified, only operators whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If \texttt{+} is appended to the command name, additional information about each operator is shown, currently just the name of the underlying function.

\dO[\+][ ]{ pattern }  
Lists collations. If \texttt{pattern} is specified, only collations whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If \texttt{+} is appended to the command name, each collation is listed with its associated description, if any. Note that only collations usable with the current database’s encoding are shown, so the results may vary in different databases of the same installation.

\dp[ ]{ pattern }  
Lists tables, views and sequences with their associated access privileges. If \texttt{pattern} is specified, only tables, views and sequences whose names match the pattern are listed.

The \texttt{GRANT} and \texttt{REVOKE} commands are used to set access privileges. The meaning of the privilege display is explained under \texttt{GRANT}.

\drds[ ]{ role-pattern [ database-pattern ] }  
Lists defined configuration settings. These settings can be role-specific, database-specific, or both. \texttt{role-pattern} and \texttt{database-pattern} are used to select specific roles and databases to list, respectively. If omitted, or if \texttt{*} is specified, all settings are listed, including those not role-specific or database-specific, respectively.

The \texttt{ALTER ROLE} and \texttt{ALTER DATABASE} commands are used to define per-role and per-database configuration settings.

\dRp[\+][ ]{ pattern }  
Lists replication publications. If \texttt{pattern} is specified, only those publications whose names match the pattern are listed. If \texttt{+} is appended to the command name, the tables associated with each publication are shown as well.

\dRs[\+][ ]{ pattern }  
Lists replication subscriptions. If \texttt{pattern} is specified, only those subscriptions whose names match the pattern are listed. If \texttt{+} is appended to the command name, additional properties of the subscriptions are shown.
\dT[S+] [ pattern ]

Lists data types. If pattern is specified, only types whose names match the pattern are listed. If + is appended to the command name, each type is listed with its internal name and size, its allowed values if it is an enum type, and its associated permissions. By default, only user-created objects are shown; supply a pattern or the S modifier to include system objects.

\du[S+] [ pattern ]

Lists database roles. (Since the concepts of “users” and “groups” have been unified into “roles”, this command is now equivalent to \dg.) By default, only user-created roles are shown; supply the S modifier to include system roles. If pattern is specified, only those roles whose names match the pattern are listed. If the form \du+ is used, additional information is shown about each role; currently this adds the comment for each role.

\dx[+] [ pattern ]

Lists installed extensions. If pattern is specified, only those extensions whose names match the pattern are listed. If the form \dx+ is used, all the objects belonging to each matching extension are listed.

\dy[+] [ pattern ]

Lists event triggers. If pattern is specified, only those event triggers whose names match the pattern are listed. If + is appended to the command name, each object is listed with its associated description.

\e or \edit [ filename ] [ line_number ]

If filename is specified, the file is edited; after the editor exits, the file's content is copied into the current query buffer. If no filename is given, the current query buffer is copied to a temporary file which is then edited in the same fashion. Or, if the current query buffer is empty, the most recently executed query is copied to a temporary file and edited in the same fashion.

The new contents of the query buffer are then re-parsed according to the normal rules of psql, treating the whole buffer as a single line. Any complete queries are immediately executed; that is, if the query buffer contains or ends with a semicolon, everything up to that point is executed. Whatever remains will wait in the query buffer; type semicolon or \s to send it, or \r to cancel it by clearing the query buffer. Treating the buffer as a single line primarily affects meta-commands: whatever is in the buffer after a meta-command will be taken as argument(s) to the meta-command, even if it spans multiple lines. (Thus you cannot make meta-command-using scripts this way. Use \i for that.)

If a line number is specified, psql will position the cursor on the specified line of the file or query buffer. Note that if a single all-digits argument is given, psql assumes it is a line number, not a file name.

\echo text [ ... ]

Prints the arguments to the standard output, separated by one space and followed by a newline. This can be useful to intersperse information in the output of scripts. For example:

```
=> \echo `date`
```
If the first argument is an unquoted -n the trailing newline is not written.

```
Tip
If you use the \o command to redirect your query output you might wish to use \qecho instead of this command.
```

```
\ef [ function_description [ line_number ] ]
```
This command fetches and edits the definition of the named function or procedure, in the form of a CREATE OR REPLACE FUNCTION or CREATE OR REPLACE PROCEDURE command. Editing is done in the same way as for \edit. After the editor exits, the updated command waits in the query buffer; type semicolon or \g to send it, or \r to cancel.

The target function can be specified by name alone, or by name and arguments, for example foo(integer, text). The argument types must be given if there is more than one function of the same name.

If no function is specified, a blank CREATE FUNCTION template is presented for editing.

If a line number is specified, psql will position the cursor on the specified line of the function body. (Note that the function body typically does not begin on the first line of the file.)

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \ef, and neither variable interpolation nor backquote expansion are performed in the arguments.

```
Tip
See under Environment for how to configure and customize your editor.
```

```
\encoding [ encoding ]
```
Sets the client character set encoding. Without an argument, this command shows the current encoding.

```
\errverbose
```
Repeats the most recent server error message at maximum verbosity, as though VERBOSITY were set to verbose and SHOW_CONTEXT were set to always.

```
\ev [ view_name [ line_number ] ]
```
This command fetches and edits the definition of the named view, in the form of a CREATE OR REPLACE VIEW command. Editing is done in the same way as for \edit. After the editor exits, the updated command waits in the query buffer; type semicolon or \g to send it, or \r to cancel.

If no view is specified, a blank CREATE VIEW template is presented for editing.

If a line number is specified, psql will position the cursor on the specified line of the view definition.

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \ev, and neither variable interpolation nor backquote expansion are performed in the arguments.
\f [ string ]

Sets the field separator for unaligned query output. The default is the vertical bar (|). It is equivalent to \pset fieldsep.

\g [ filename ]
\g [ command ]

Sends the current query buffer to the server for execution. If an argument is given, the query's output is written to the named file or piped to the given shell command, instead of displaying it as usual. The file or command is written to only if the query successfully returns zero or more tuples, not if the query fails or is a non-data-returning SQL command.

If the current query buffer is empty, the most recently sent query is re-executed instead. Except for that behavior, \g without an argument is essentially equivalent to a semicolon. A \g with argument is a "one-shot" alternative to the \o command.

If the argument begins with |, then the entire remainder of the line is taken to be the command to execute, and neither variable interpolation nor backquote expansion are performed in it. The rest of the line is simply passed literally to the shell.

\gdesc

Shows the description (that is, the column names and data types) of the result of the current query buffer. The query is not actually executed; however, if it contains some type of syntax error, that error will be reported in the normal way.

If the current query buffer is empty, the most recently sent query is described instead.

\gexec

Sends the current query buffer to the server, then treats each column of each row of the query's output (if any) as a SQL statement to be executed. For example, to create an index on each column of my_table:

```bash
=> SELECT format('create index on my_table(%I)', attname) -> FROM pg_attribute
  -> WHERE attrelid = 'my_table'::regclass AND attnum > 0
  -> ORDER BY attnum
  -> \gexec
CREATE INDEX
CREATE INDEX
CREATE INDEX
CREATE INDEX
```

The generated queries are executed in the order in which the rows are returned, and left-to-right within each row if there is more than one column. NULL fields are ignored. The generated queries are sent literally to the server for processing, so they cannot be psql meta-commands nor contain psql variable references. If any individual query fails, execution of the remaining queries continues unless ON_ERROR_STOP is set. Execution of each query is subject to ECHO processing. (Setting ECHO to all or queries is often advisable when using \gexec.) Query logging, single-step mode, timing, and other query execution features apply to each generated query as well.

If the current query buffer is empty, the most recently sent query is re-executed instead.

\gset [ prefix ]

Sends the current query buffer to the server and stores the query's output into psql variables (see Variables). The query to be executed must return exactly one row. Each column of the row is stored into a separate variable, named the same as the column. For example:
=> SELECT 'hello' AS var1, 10 AS var2
-> \gset
=> \echo :var1 :var2
hello 10

If you specify a **prefix**, that string is prepended to the query's column names to create the variable names to use:

=> SELECT 'hello' AS var1, 10 AS var2
-> \gset result_
=> \echo :result_var1 :result_var2
hello 10

If a column result is NULL, the corresponding variable is unset rather than being set.

If the query fails or does not return one row, no variables are changed.

If the current query buffer is empty, the most recently sent query is re-executed instead.

\gx [ filename ]
\gx [ | command ]

\gx is equivalent to \g, but forces expanded output mode for this query. See \x.

\h or \help [ command ]

Gives syntax help on the specified SQL command. If *command* is not specified, then psql will list all the commands for which syntax help is available. If *command* is an asterisk (*), then syntax help on all SQL commands is shown.

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \help, and neither variable interpolation nor backquote expansion are performed in the arguments.

**Note**

To simplify typing, commands that consists of several words do not have to be quoted. Thus it is fine to type \help alter table.

\H or \html

Turns on HTML query output format. If the HTML format is already on, it is switched back to the default aligned text format. This command is for compatibility and convenience, but see \pset about setting other output options.

\i or \include filename

Reads input from the file *filename* and executes it as though it had been typed on the keyboard.

If *filename* is - (hyphen), then standard input is read until an EOF indication or \q meta-command. This can be used to intersperse interactive input with input from files. Note that Readline behavior will be used only if it is active at the outermost level.

**Note**

If you want to see the lines on the screen as they are read you must set the variable ECHO to all.

1914
This group of commands implements nestable conditional blocks. A conditional block must begin with an `\if` and end with an `\endif`. In between there may be any number of `\elif` clauses, which may optionally be followed by a single `\else` clause. Ordinary queries and other types of backslash commands may (and usually do) appear between the commands forming a conditional block.

The `\if` and `\elif` commands read their argument(s) and evaluate them as a boolean expression. If the expression yields `true` then processing continues normally; otherwise, lines are skipped until a matching `\elif`, `\else`, or `\endif` is reached. Once an `\if` or `\elif` test has succeeded, the arguments of later `\elif` commands in the same block are not evaluated but are treated as false. Lines following an `\else` are processed only if no earlier matching `\if` or `\elif` succeeded.

The `expression` argument of an `\if` or `\elif` command is subject to variable interpolation and backquote expansion, just like any other backslash command argument. After that it is evaluated like the value of an on/off option variable. So a valid value is any unambiguous case-insensitive match for one of: `true`, `false`, `1`, `0`, `on`, `off`, `yes`, `no`. For example, `t`, `T`, and `tR` will all be considered to be `true`.

Expressions that do not properly evaluate to true or false will generate a warning and be treated as false.

Lines being skipped are parsed normally to identify queries and backslash commands, but queries are not sent to the server, and backslash commands other than conditionals (`\if`, `\elif`, `\else`, `\endif`) are ignored. Conditional commands are checked only for valid nesting. Variable references in skipped lines are not expanded, and backquote expansion is not performed either.

All the backslash commands of a given conditional block must appear in the same source file. If EOF is reached on the main input file or an `\include`-ed file before all local `\if`-blocks have been closed, then `psql` will raise an error.

Here is an example:

```sql
-- check for the existence of two separate records in the database and store
-- the results in separate psql variables
SELECT
  EXISTS(SELECT 1 FROM customer WHERE customer_id = 123) as is_customer,
  EXISTS(SELECT 1 FROM employee WHERE employee_id = 456) as is_employee
\gset
\if :is_customer
  SELECT * FROM customer WHERE customer_id = 123;
\elif :is_employee
  \echo 'is not a customer but is an employee'
  SELECT * FROM employee WHERE employee_id = 456;
\else
  \if yes
    \echo 'not a customer or employee'
  \else
    \echo 'this will never print'
  \endif
\endif
```

\end{verbatim}

The \texttt{\ir} command is similar to \texttt{i}, but resolves relative file names differently. When executing in interactive mode, the two commands behave identically. However, when invoked from a script, \texttt{\ir} interprets file names relative to the directory in which the script is located, rather than the current working directory.

\texttt{\l[+]} or \texttt{\list[+] [ pattern ]}

List the databases in the server and show their names, owners, character set encodings, and access privileges. If \texttt{pattern} is specified, only databases whose names match the pattern are listed. If + is appended to the command name, database sizes, default tablespaces, and descriptions are also displayed. (Size information is only available for databases that the current user can connect to.)

\texttt{\lo_export loid filename}

Reads the large object with OID \texttt{loid} from the database and writes it to \texttt{filename}. Note that this is subtly different from the server function \texttt{lo_export}, which acts with the permissions of the user that the database server runs as and on the server's file system.

\begin{verbatim}
Tip
Use \texttt{\lo_list} to find out the large object's OID.
\end{verbatim}

\texttt{\lo_import filename [ comment ]}

Stores the file into a PostgreSQL large object. Optionally, it associates the given comment with the object. Example:

\begin{verbatim}
foo=> \texttt{\lo_import '/home/peter/pictures/photo.xcf' 'a picture of me'}
\lo_import 152801
\end{verbatim}

The response indicates that the large object received object ID 152801, which can be used to access the newly-created large object in the future. For the sake of readability, it is recommended to always associate a human-readable comment with every object. Both OIDs and comments can be viewed with the \texttt{\lo_list} command.

Note that this command is subtly different from the server-side \texttt{lo_import} because it acts as the local user on the local file system, rather than the server's user and file system.

\texttt{\lo_list}

Shows a list of all PostgreSQL large objects currently stored in the database, along with any comments provided for them.

\texttt{\lo_unlink loid}

Deletes the large object with OID \texttt{loid} from the database.

\begin{verbatim}
Tip
Use \texttt{\lo_list} to find out the large object's OID.
\end{verbatim}
\o or \out [ filename ]
\o or \out [ command ]

Arranges to save future query results to the file filename or pipe future results to the shell command command. If no argument is specified, the query output is reset to the standard output.

If the argument begins with |, then the entire remainder of the line is taken to be the command to execute, and neither variable interpolation nor backquote expansion are performed in it. The rest of the line is simply passed literally to the shell.

“Query results” includes all tables, command responses, and notices obtained from the database server, as well as output of various backslash commands that query the database (such as \d); but not error messages.

Tip
To intersperse text output in between query results, use \qecho.

\p or \print
Print the current query buffer to the standard output. If the current query buffer is empty, the most recently executed query is printed instead.

\password [ username ]
Changes the password of the specified user (by default, the current user). This command prompts for the new password, encrypts it, and sends it to the server as an ALTER ROLE command. This makes sure that the new password does not appear in cleartext in the command history, the server log, or elsewhere.

\prompt [ text ] name
Prompts the user to supply text, which is assigned to the variable name. An optional prompt string, text, can be specified. (For multiword prompts, surround the text with single quotes.)

By default, \prompt uses the terminal for input and output. However, if the \-f command line switch was used, \prompt uses standard input and standard output.

\pset [ option [ value ] ]
This command sets options affecting the output of query result tables. option indicates which option is to be set. The semantics of value vary depending on the selected option. For some options, omitting value causes the option to be toggled or unset, as described under the particular option. If no such behavior is mentioned, then omitting value just results in the current setting being displayed.

\pset without any arguments displays the current status of all printing options.

Adjustable printing options are:

border

The value must be a number. In general, the higher the number the more borders and lines the tables will have, but details depend on the particular format. In HTML format, this will translate directly into the border=... attribute. In most other formats only values 0 (no border), 1 (internal dividing lines), and 2 (table frame) make sense, and values above 2 will be treated the same as border = 2. The latex and latex-longtable formats additionally allow a value of 3 to add dividing lines between data rows.
columns

Sets the target width for the wrapped format, and also the width limit for determining whether output is wide enough to require the pager or switch to the vertical display in expanded auto mode. Zero (the default) causes the target width to be controlled by the environment variable COLUMNS, or the detected screen width if COLUMNS is not set. In addition, if columns is zero then the wrapped format only affects screen output. If columns is nonzero then file and pipe output is wrapped to that width as well.

expanded (or x)

If value is specified it must be either on or off, which will enable or disable expanded mode, or auto. If value is omitted the command toggles between the on and off settings. When expanded mode is enabled, query results are displayed in two columns, with the column name on the left and the data on the right. This mode is useful if the data wouldn’t fit on the screen in the normal “horizontal” mode. In the auto setting, the expanded mode is used whenever the query output has more than one column and is wider than the screen; otherwise, the regular mode is used. The auto setting is only effective in the aligned and wrapped formats. In other formats, it always behaves as if the expanded mode is off.

fieldsep

Specifies the field separator to be used in unaligned output format. That way one can create, for example, tab- or comma-separated output, which other programs might prefer. To set a tab as field separator, type \pset fieldsep \t. The default field separator is '|' (a vertical bar).

fieldsep_zero

Sets the field separator to use in unaligned output format to a zero byte.

footer

If value is specified it must be either on or off which will enable or disable display of the table footer (the \( n \) rows count). If value is omitted the command toggles footer display on or off.

format

Sets the output format to one of unaligned, aligned, wrapped, html, asciidoc, latex (uses tabular), latex-longtable, or troff-ms. Unique abbreviations are allowed. (That would mean one letter is enough.)

unaligned format writes all columns of a row on one line, separated by the currently active field separator. This is useful for creating output that might be intended to be read in by other programs (for example, tab-separated or comma-separated format).

aligned format is the standard, human-readable, nicely formatted text output; this is the default.

wrapped format is like aligned but wraps wide data values across lines to make the output fit in the target column width. The target width is determined as described under the columns option. Note that psql will not attempt to wrap column header titles; therefore, wrapped format behaves the same as aligned if the total width needed for column headers exceeds the target.

The html, asciidoc, latex, latex-longtable, and troff-ms formats put out tables that are intended to be included in documents using the respective mark-up language. They are not complete documents! This might not be necessary in HTML, but in LaTeX you must have a complete document wrapper. latex-longtable also requires the LaTeX longtable and booktabs packages.
linestyle

Sets the border line drawing style to one of ascii, old-ascii, or unicode. Unique abbreviations are allowed. (That would mean one letter is enough.) The default setting is ascii. This option only affects the aligned and wrapped output formats.

ascii style uses plain ASCII characters. Newlines in data are shown using a + symbol in the right-hand margin. When the wrapped format wraps data from one line to the next without a newline character, a dot (.) is shown in the right-hand margin of the first line, and again in the left-hand margin of the following line.

old-ascii style uses plain ASCII characters, using the formatting style used in PostgreSQL 8.4 and earlier. Newlines in data are shown using a ; symbol in place of the left-hand column separator. When the data is wrapped from one line to the next without a newline character, a ; symbol is used in place of the left-hand column separator.

unicode style uses Unicode box-drawing characters. Newlines in data are shown using a carriage return symbol in the right-hand margin. When the data is wrapped from one line to the next without a newline character, an ellipsis symbol is shown in the right-hand margin of the first line, and again in the left-hand margin of the following line.

When the border setting is greater than zero, the linestyle option also determines the characters with which the border lines are drawn. Plain ASCII characters work everywhere, but Unicode characters look nicer on displays that recognize them.

null

Sets the string to be printed in place of a null value. The default is to print nothing, which can easily be mistaken for an empty string. For example, one might prefer \pset null '(null)'.

numericlocale

If value is specified it must be either on or off which will enable or disable display of a locale-specific character to separate groups of digits to the left of the decimal marker. If value is omitted the command toggles between regular and locale-specific numeric output.

pager

Controls use of a pager program for query and psql help output. If the environment variable PSQL_PAGER or PAGER is set, the output is piped to the specified program. Otherwise a platform-dependent default program (such as more) is used.

When the pager option is off, the pager program is not used. When the pager option is on, the pager is used when appropriate, i.e., when the output is to a terminal and will not fit on the screen. The pager option can also be set to always, which causes the pager to be used for all terminal output regardless of whether it fits on the screen. \pset pager without a value toggles pager use on and off.

pager_min_lines

If pager_min_lines is set to a number greater than the page height, the pager program will not be called unless there are at least this many lines of output to show. The default setting is 0.

recordsep

Specifies the record (line) separator to use in unaligned output format. The default is a newline character.

recordsep_zero

Sets the record separator to use in unaligned output format to a zero byte.
tableattr (or T)

In HTML format, this specifies attributes to be placed inside the table tag. This could for example be cellpadding or bgcolor. Note that you probably don't want to specify border here, as that is already taken care of by \pset border. If no \value is given, the table attributes are unset.

In latex-longtable format, this controls the proportional width of each column containing a left-aligned data type. It is specified as a whitespace-separated list of values, e.g. '0.2 0.2 0.6'. Unspecified output columns use the last specified value.

title (or C)

Sets the table title for any subsequently printed tables. This can be used to give your output descriptive tags. If no \value is given, the title is unset.

tuples_only (or t)

If \value is specified it must be either on or off which will enable or disable tuples-only mode. If \value is omitted the command toggles between regular and tuples-only output. Regular output includes extra information such as column headers, titles, and various footers. In tuples-only mode, only actual table data is shown.

unicode_border_linestyle

Sets the border drawing style for the unicode line style to one of single or double.

unicode_column_linestyle

Sets the column drawing style for the unicode line style to one of single or double.

unicode_header_linestyle

Sets the header drawing style for the unicode line style to one of single or double.

Illustrations of how these different formats look can be seen in the Examples section.

Tip

There are various shortcut commands for \pset. See \a, \C, \f, \H, \t, \T, and \x.

\q or \quit

Quits the psql program. In a script file, only execution of that script is terminated.

\gecho text [ ... ]

This command is identical to \echo except that the output will be written to the query output channel, as set by \o.

\r or \reset

Resets (clears) the query buffer.

\s [ filename ]

Print psql's command line history to filename. If filename is omitted, the history is written to the standard output (using the pager if appropriate). This command is not available if psql was built without Readline support.
\set [ name [ value [ ... ] ] ]

Sets the psql variable name to value, or if more than one value is given, to the concatenation of all of them. If only one argument is given, the variable is set to an empty-string value. To unset a variable, use the \unset command.

\set without any arguments displays the names and values of all currently-set psql variables.

Valid variable names can contain letters, digits, and underscores. See the section Variables below for details. Variable names are case-sensitive.

Certain variables are special, in that they control psql's behavior or are automatically set to reflect connection state. These variables are documented in Variables, below.

\setenv name [ value ]

Sets the environment variable name to value, or if the value is not supplied, unsets the environment variable. Example:

```
  testdb=> \setenv PAGER less
  testdb=> \setenv LESS -imx4F
```

\sf[++]

This command fetches and shows the definition of the named function or procedure, in the form of a CREATE OR REPLACE FUNCTION or CREATE OR REPLACE PROCEDURE command. The definition is printed to the current query output channel, as set by \o.

The target function can be specified by name alone, or by name and arguments, for example foo(integer, text). The argument types must be given if there is more than one function of the same name.

If + is appended to the command name, then the output lines are numbered, with the first line of the function body being line 1.

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \sf, and neither variable interpolation nor backquote expansion are performed in the arguments.

\sv[++]

This command fetches and shows the definition of the named view, in the form of a CREATE OR REPLACE VIEW command. The definition is printed to the current query output channel, as set by \o.

If + is appended to the command name, then the output lines are numbered from 1.

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \sv, and neither variable interpolation nor backquote expansion are performed in the arguments.

\t

Toggles the display of output column name headings and row count footer. This command is equivalent to \pset tuples_only and is provided for convenience.
\T table_options

Specifies attributes to be placed within the table tag in HTML output format. This command is equivalent to \pset tableattr table_options.

\timing [ on | off ]

With a parameter, turns displaying of how long each SQL statement takes on or off. Without a parameter, toggles the display between on and off. The display is in milliseconds; intervals longer than 1 second are also shown in minutes:seconds format, with hours and days fields added if needed.

\unset name

Unsets (deletes) the psql variable name.

Most variables that control psql’s behavior cannot be unset; instead, an \unset command is interpreted as setting them to their default values. See Variables, below.

\w or \write filename
\w or \write | command

Writes the current query buffer to the file filename or pipes it to the shell command command. If the current query buffer is empty, the most recently executed query is written instead.

If the argument begins with |, then the entire remainder of the line is taken to be the command to execute, and neither variable interpolation nor backquote expansion are performed in it. The rest of the line is simply passed literally to the shell.

\watch [ seconds ]

Repeatedly execute the current query buffer (as \g does) until interrupted or the query fails. Wait the specified number of seconds (default 2) between executions. Each query result is displayed with a header that includes the \pset title string (if any), the time as of query start, and the delay interval.

If the current query buffer is empty, the most recently sent query is re-executed instead.

\x [ on | off | auto ]

Sets or toggles expanded table formatting mode. As such it is equivalent to \pset expanded.

\z [ pattern ]

Lists tables, views and sequences with their associated access privileges. If a pattern is specified, only tables, views and sequences whose names match the pattern are listed.

This is an alias for \dp (“display privileges”).

\! [ command ]

With no argument, escapes to a sub-shell; psql resumes when the sub-shell exits. With an argument, executes the shell command command.

Unlike most other meta-commands, the entire remainder of the line is always taken to be the argument(s) of \!, and neither variable interpolation nor backquote expansion are performed in the arguments. The rest of the line is simply passed literally to the shell.

\? [ topic ]

Shows help information. The optional topic parameter (defaulting to commands) selects which part of psql is explained: commands describes psql’s backslash commands; options describes the command-line options that can be passed to psql; and variables shows help about psql configuration variables.
Backslash-semicolons are not meta-commands in the same way as the preceding commands; rather, it simply causes a semicolon to be added to the query buffer without any further processing.

Normally, psql will dispatch a SQL command to the server as soon as it reaches the command-ending semicolon, even if more input remains on the current line. Thus for example entering

```
select 1; select 2; select 3;
```

will result in the three SQL commands being individually sent to the server, with each one's results being displayed before continuing to the next command. However, a semicolon entered as `\;` will not trigger command processing, so that the command before it and the one after are effectively combined and sent to the server in one request. So for example

```
select 1\; select 2\; select 3;
```

results in sending the three SQL commands to the server in a single request, when the non-backslashed semicolon is reached. The server executes such a request as a single transaction, unless there are explicit BEGIN/COMMIT commands included in the string to divide it into multiple transactions. (See Section 53.2.2.1 for more details about how the server handles multi-query strings.) psql prints only the last query result it receives for each request; in this example, although all three `SELECT` statements are indeed executed, psql only prints the 3.

**Patterns**

The various \d commands accept a **pattern** parameter to specify the object name(s) to be displayed. In the simplest case, a pattern is just the exact name of the object. The characters within a pattern are normally folded to lower case, just as in SQL names; for example, `\dt 'FOO'` will display the table named `foo`. As in SQL names, placing double quotes around a pattern stops folding to lower case. Should you need to include an actual double quote character in a pattern, write it as a pair of double quotes within a double-quote sequence; again this is in accord with the rules for SQL quoted identifiers. For example, `\dt "FOO""BAR"` will display the table named `FOO"BAR` (not `foo"bar`). Unlike the normal rules for SQL names, you can put double quotes around just part of a pattern, for instance `\dt 'FOO"BAR'` will display the table named `fooFOObar`.

Whenever the **pattern** parameter is omitted completely, the \d commands display all objects that are visible in the current schema search path — this is equivalent to using `*` as the pattern. (An object is said to be **visible** if its containing schema is in the search path and no object of the same kind and name appears earlier in the search path. This is equivalent to the statement that the object can be referenced by name without explicit schema qualification.) To see all objects in the database regardless of visibility, use `*.*` as the pattern.

Within a pattern, `*` matches any sequence of characters (including no characters) and `?` matches any single character. (This notation is comparable to Unix shell file name patterns.) For example, `\dt 'int*'` displays tables whose names begin with `int`. But within double quotes, `*` and `?` lose these special meanings and are just matched literally.

A pattern that contains a dot (`.`) is interpreted as a schema name pattern followed by an object name pattern. For example, `\dt 'foo*.bar'` displays all tables whose table name includes `bar` that are in schemas whose schema name starts with `foo`. When no dot appears, then the pattern matches only objects that are visible in the current schema search path. Again, a dot within double quotes loses its special meaning and is matched literally.

Advanced users can use regular-expression notations such as character classes, for example `[0-9]` to match any digit. All regular expression special characters work as specified in Section 9.7.3, except for `.` which is taken as a separator as mentioned above, `*` which is translated to the regular-expression notation `.*`, `?` which is translated to `.*`, and `\` which is matched literally. You can emulate these
pattern characters at need by writing \? for ., \( R+ \mid \) for \( R^* \), or \( R \mid \) for \( R？ \). $ is not needed as a regular-expression character since the pattern must match the whole name, unlike the usual interpretation of regular expressions (in other words, $ is automatically appended to your pattern). Write * at the beginning and/or end if you don’t wish the pattern to be anchored. Note that within double quotes, all regular expression special characters lose their special meanings and are matched literally. Also, the regular expression special characters are matched literally in operator name patterns (i.e., the argument of \do).

**Advanced Features**

**Variables**

psql provides variable substitution features similar to common Unix command shells. Variables are simply name/value pairs, where the value can be any string of any length. The name must consist of letters (including non-Latin letters), digits, and underscores.

To set a variable, use the psql meta-command \set. For example,

```

testdb=> \set foo bar
```

sets the variable foo to the value bar. To retrieve the content of the variable, precede the name with a colon, for example:

```

testdb=> \echo :foo
```

bar

This works in both regular SQL commands and meta-commands; there is more detail in SQL Interpolation, below.

If you call \set without a second argument, the variable is set to an empty-string value. To unset (i.e., delete) a variable, use the command \unset. To show the values of all variables, call \set without any argument.

```

<table>
<thead>
<tr>
<th>Note</th>
</tr>
</thead>
</table>

The arguments of \set are subject to the same substitution rules as with other commands. Thus you can construct interesting references such as \set :foo ‘something’ and get “soft links” or “variable variables” of Perl or PHP fame, respectively. Unfortunately (or fortunately?), there is no way to do anything useful with these constructs. On the other hand, \set bar :foo is a perfectly valid way to copy a variable.

A number of these variables are treated specially by psql. They represent certain option settings that can be changed at run time by altering the value of the variable, or in some cases represent changeable state of psql. By convention, all specially treated variables’ names consist of all upper-case ASCII letters (and possibly digits and underscores). To ensure maximum compatibility in the future, avoid using such variable names for your own purposes.

Variables that control psql’s behavior generally cannot be unset or set to invalid values. An \unset command is allowed but is interpreted as setting the variable to its default value. A \set command without a second argument is interpreted as setting the variable to on, for control variables that accept that value, and is rejected for others. Also, control variables that accept the values on and off will also accept other common spellings of Boolean values, such as true and false.

The specially treated variables are:
AUTOCOMMIT

When on (the default), each SQL command is automatically committed upon successful completion. To postpone commit in this mode, you must enter a BEGIN or START TRANSACTION SQL command. When off or unset, SQL commands are not committed until you explicitly issue COMMIT or END. The autocommit-off mode works by issuing an implicit BEGIN for you, just before any command that is not already in a transaction block and is not itself a BEGIN or other transaction-control command, nor a command that cannot be executed inside a transaction block (such as VACUUM).

Note

In autocommit-off mode, you must explicitly abandon any failed transaction by entering ABORT or ROLLBACK. Also keep in mind that if you exit the session without committing, your work will be lost.

Note

The autocommit-on mode is PostgreSQL's traditional behavior, but autocommit-off is closer to the SQL spec. If you prefer autocommit-off, you might wish to set it in the system-wide psqlrc file or your ~/.psqlrc file.

COMP_KEYWORD_CASE

Determines which letter case to use when completing an SQL key word. If set to lower or upper, the completed word will be in lower or upper case, respectively. If set to preserve-lower or preserve-upper (the default), the completed word will be in the case of the word already entered, but words being completed without anything entered will be in lower or upper case, respectively.

DBNAME

The name of the database you are currently connected to. This is set every time you connect to a database (including program start-up), but can be changed or unset.

ECHO

If set to all, all nonempty input lines are printed to standard output as they are read. (This does not apply to lines read interactively.) To select this behavior on program start-up, use the switch -a. If set to queries, psql prints each query to standard output as it is sent to the server. The switch to select this behavior is -e. If set to errors, then only failed queries are displayed on standard error output. The switch for this behavior is -b. If set to none (the default), then no queries are displayed.

ECHO_HIDDEN

When this variable is set to on and a backslash command queries the database, the query is first shown. This feature helps you to study PostgreSQL internals and provide similar functionality in your own programs. (To select this behavior on program start-up, use the switch -E.) If you set this variable to the value noexec, the queries are just shown but are not actually sent to the server and executed. The default value is off.

ENCODING

The current client character set encoding. This is set every time you connect to a database (including program start-up), and when you change the encoding with \encoding, but it can be changed or unset.
**ERROR**

true if the last SQL query failed, false if it succeeded. See also SQLSTATE.

**FETCH_COUNT**

If this variable is set to an integer value greater than zero, the results of SELECT queries are fetched and displayed in groups of that many rows, rather than the default behavior of collecting the entire result set before display. Therefore only a limited amount of memory is used, regardless of the size of the result set. Settings of 100 to 1000 are commonly used when enabling this feature. Keep in mind that when using this feature, a query might fail after having already displayed some rows.

**Tip**

Although you can use any output format with this feature, the default aligned format tends to look bad because each group of FETCH_COUNT rows will be formatted separately, leading to varying column widths across the row groups. The other output formats work better.

**HISTCONTROL**

If this variable is set to ignore space, lines which begin with a space are not entered into the history list. If set to a value of ignoredups, lines matching the previous history line are not entered. A value of ignoreboth combines the two options. If set to none (the default), all lines read in interactive mode are saved on the history list.

**Note**

This feature was shamelessly plagiarized from Bash.

**HISTFILE**

The file name that will be used to store the history list. If unset, the file name is taken from the PSQL_HISTORY environment variable. If that is not set either, the default is ~/.psql_history, or %APPDATA%/postgresql/psql_history on Windows. For example, putting:

\set HISTFILE ~/.psql_history- :DBNAME

in ~/.psqlrc will cause psql to maintain a separate history for each database.

**Note**

This feature was shamelessly plagiarized from Bash.

**HISTSIZE**

The maximum number of commands to store in the command history (default 500). If set to a negative value, no limit is applied.

**Note**

This feature was shamelessly plagiarized from Bash.
HOST

The database server host you are currently connected to. This is set every time you connect to a database (including program start-up), but can be changed or unset.

IGNOREEOF

If set to 1 or less, sending an EOF character (usually Control-D) to an interactive session of psql will terminate the application. If set to a larger numeric value, that many consecutive EOF characters must be typed to make an interactive session terminate. If the variable is set to a non-numeric value, it is interpreted as 10. The default is 0.

Note

This feature was shamelessly plagiarized from Bash.

LASTOID

The value of the last affected OID, as returned from an INSERT or \lo_import command. This variable is only guaranteed to be valid until after the result of the next SQL command has been displayed.

LAST_ERROR_MESSAGE
LAST_ERROR_SQLSTATE

The primary error message and associated SQLSTATE code for the most recent failed query in the current psql session, or an empty string and 00000 if no error has occurred in the current session.

ON_ERROR_ROLLBACK

When set to on, if a statement in a transaction block generates an error, the error is ignored and the transaction continues. When set to interactive, such errors are only ignored in interactive sessions, and not when reading script files. When set to off (the default), a statement in a transaction block that generates an error aborts the entire transaction. The error rollback mode works by issuing an implicit SAVEPOINT for you, just before each command that is in a transaction block, and then rolling back to the savepoint if the command fails.

ON_ERROR_STOP

By default, command processing continues after an error. When this variable is set to on, processing will instead stop immediately. In interactive mode, psql will return to the command prompt; otherwise, psql will exit, returning error code 3 to distinguish this case from fatal error conditions, which are reported using error code 1. In either case, any currently running scripts (the top-level script, if any, and any other scripts which it may have in invoked) will be terminated immediately. If the top-level command string contained multiple SQL commands, processing will stop with the current command.

PORT

The database server port to which you are currently connected. This is set every time you connect to a database (including program start-up), but can be changed or unset.

PROMPT1
PROMPT2
PROMPT3

These specify what the prompts psql issues should look like. See Prompting below.
QUIET

Setting this variable to on is equivalent to the command line option -q. It is probably not too useful in interactive mode.

ROW_COUNT

The number of rows returned or affected by the last SQL query, or 0 if the query failed or did not report a row count.

SERVER_VERSION_NAME
SERVER_VERSION_NUM

The server's version number as a string, for example 9.6.2, 10.1 or 11beta1, and in numeric form, for example 90602 or 100001. These are set every time you connect to a database (including program start-up), but can be changed or unset.

SHOW_CONTEXT

This variable can be set to the values never, errors, or always to control whether CONTEXT fields are displayed in messages from the server. The default is errors (meaning that context will be shown in error messages, but not in notice or warning messages). This setting has no effect when VERBOSITY is set to terse. (See also \
errverbose, for use when you want a verbose version of the error you just got.)

SINGLELINE

Setting this variable to on is equivalent to the command line option -S.

SINGLESTEP

Setting this variable to on is equivalent to the command line option -s.

SQLSTATE

The error code (see Appendix A) associated with the last SQL query's failure, or 00000 if it succeeded.

USER

The database user you are currently connected as. This is set every time you connect to a database (including program start-up), but can be changed or unset.

VERBOSITY

This variable can be set to the values default, verbose, or terse to control the verbosity of error reports. (See also \errverbose, for use when you want a verbose version of the error you just got.)

VERSION
VERSION_NAME
VERSION_NUM

These variables are set at program start-up to reflect psql's version, respectively as a verbose string, a short string (e.g., 9.6.2, 10.1, or 11beta1), and a number (e.g., 90602 or 100001). They can be changed or unset.

**SQL Interpolation**

A key feature of psql variables is that you can substitute ("interpolate") them into regular SQL statements, as well as the arguments of meta-commands. Furthermore, psql provides facilities for ensuring
that variable values used as SQL literals and identifiers are properly quoted. The syntax for interpolating a value without any quoting is to prepend the variable name with a colon (:). For example,

```
testdb=> \set foo 'my_table'
testdb=> SELECT * FROM :foo;
```

would query the table my_table. Note that this may be unsafe: the value of the variable is copied literally, so it can contain unbalanced quotes, or even backslash commands. You must make sure that it makes sense where you put it.

When a value is to be used as an SQL literal or identifier, it is safest to arrange for it to be quoted. To quote the value of a variable as an SQL literal, write a colon followed by the variable name in single quotes. To quote the value as an SQL identifier, write a colon followed by the variable name in double quotes. These constructs deal correctly with quotes and other special characters embedded within the variable value. The previous example would be more safely written this way:

```
testdb=> \set foo 'my_table'
testdb=> SELECT * FROM :"foo";
```

Variable interpolation will not be performed within quoted SQL literals and identifiers. Therefore, a construction such as ':foo' doesn't work to produce a quoted literal from a variable's value (and it would be unsafe if it did work, since it wouldn't correctly handle quotes embedded in the value).

One example use of this mechanism is to copy the contents of a file into a table column. First load the file into a variable and then interpolate the variable's value as a quoted string:

```
testdb=> \set content `cat my_file.txt`
testdb=> INSERT INTO my_table VALUES (:'content');
```

(Note that this still won't work if my_file.txt contains NUL bytes. psql does not support embedded NUL bytes in variable values.)

Since colons can legally appear in SQL commands, an apparent attempt at interpolation (that is, :name, :'name', or :"name") is not replaced unless the named variable is currently set. In any case, you can escape a colon with a backslash to protect it from substitution.

The :{?name} special syntax returns TRUE or FALSE depending on whether the variable exists or not, and is thus always substituted, unless the colon is backslash-escaped.

The colon syntax for variables is standard SQL for embedded query languages, such as ECPG. The colon syntaxes for array slices and type casts are PostgreSQL extensions, which can sometimes conflict with the standard usage. The colon-quote syntax for escaping a variable's value as an SQL literal or identifier is a psql extension.

**Prompts**

The prompts psql issues can be customized to your preference. The three variables PROMPT1, PROMPT2, and PROMPT3 contain strings and special escape sequences that describe the appearance of the prompt. Prompt 1 is the normal prompt that is issued when psql requests a new command. Prompt 2 is issued when more input is expected during command entry, for example because the command was not terminated with a semicolon or a quote was not closed. Prompt 3 is issued when you are running an SQL COPY FROM STDIN command and you need to type in a row value on the terminal.

The value of the selected prompt variable is printed literally, except where a percent sign (%) is encountered. Depending on the next character, certain other text is substituted instead. Defined substitutions are:
%M

The full host name (with domain name) of the database server, or [local] if the connection is over a Unix domain socket, or [local:/dir/name], if the Unix domain socket is not at the compiled in default location.

%m

The host name of the database server, truncated at the first dot, or [local] if the connection is over a Unix domain socket.

%>

The port number at which the database server is listening.

%n

The database session user name. (The expansion of this value might change during a database session as the result of the command SET SESSION AUTHORIZATION.)

%/

The name of the current database.

%~

Like %/, but the output is ~ (tilde) if the database is your default database.

%#

If the session user is a database superuser, then a #, otherwise a >. (The expansion of this value might change during a database session as the result of the command SET SESSION AUTHORIZATION.)

%p

The process ID of the backend currently connected to.

%R

In prompt 1 normally =, but @ if the session is in an inactive branch of a conditional block, or ^ if in single-line mode, or ! if the session is disconnected from the database (which can happen if \connect fails). In prompt 2 %R is replaced by a character that depends on why psql expects more input: = if the command simply wasn't terminated yet, but * if there is an unfinished / * ... */ comment, a single quote if there is an unfinished quoted string, a double quote if there is an unfinished quoted identifier, a dollar sign if there is an unfinished dollar-quoted string, or ( if there is an unmatched left parenthesis. In prompt 3 %R doesn't produce anything.

%x

Transaction status: an empty string when not in a transaction block, or * when in a transaction block, or ! when in a failed transaction block, or ? when the transaction state is indeterminate (for example, because there is no connection).

%l

The line number inside the current statement, starting from 1.

%digits

The character with the indicated octal code is substituted.
\%: \texttt{name}

The value of the \texttt{psql} variable \texttt{name}. See the section Variables for details.

\%\ `\texttt{command}\`

The output of \texttt{command}, similar to ordinary "back-tick" substitution.

\% [ ... \%]

Prompts can contain terminal control characters which, for example, change the color, background, or style of the prompt text, or change the title of the terminal window. In order for the line editing features of \texttt{Readline} to work properly, these non-printing control characters must be designated as invisible by surrounding them with \% [ and \%]. Multiple pairs of these can occur within the prompt. For example:

\begin{verbatim}
testdb=> \set PROMPT1 '%[\%033[1;33;40m%]n@%/R%[\%033[0m%]# ' 
\end{verbatim}

results in a boldfaced (1;) yellow-on-black (33;40) prompt on VT100-compatible, color-capable terminals.

To insert a percent sign into your prompt, write \%%. The default prompts are '%/R%#' for prompts 1 and 2, and '>> ' for prompt 3.

\begin{center}
\textbf{Note}
\end{center}

This feature was shamelessly plagiarized from tcsh.

\textbf{Command-Line Editing}

\texttt{psql} supports the \texttt{Readline} library for convenient line editing and retrieval. The command history is automatically saved when \texttt{psql} exits and is reloaded when \texttt{psql} starts up. Tab-completion is also supported, although the completion logic makes no claim to be an SQL parser. The queries generated by tab-completion can also interfere with other SQL commands, e.g. \texttt{SET TRANSACTION ISOLATION LEVEL}. If for some reason you do not like the tab completion, you can turn it off by putting this in a file named \texttt{.inputrc} in your home directory:

\begin{verbatim}
$if psql
set disable-completion on
$endif
\end{verbatim}

(This is not a \texttt{psql} but a \texttt{Readline} feature. Read its documentation for further details.)

\textbf{Environment}

\texttt{COLUMNS}

If \texttt{\pset columns} is zero, controls the width for the \texttt{wrapped} format and width for determining if wide output requires the pager or should be switched to the vertical format in expanded auto mode.

\texttt{PGDATABASE}
\texttt{PGHOST}
\texttt{PGPORT}
\texttt{PGUSER}

Default connection parameters (see Section 34.14).
Editor used by the \e, \ef, and \ev commands. These variables are examined in the order listed; the first that is set is used. If none of them is set, the default is to use vi on Unix systems or notepad.exe on Windows systems.

When \e, \ef, or \ev is used with a line number argument, this variable specifies the command-line argument used to pass the starting line number to the user's editor. For editors such as Emacs or vi, this is a plus sign. Include a trailing space in the value of the variable if there needs to be space between the option name and the line number. Examples:

\[ \text{PSQL_EDITOR_LINENUMBER_ARG} = '+' \]
\[ \text{PSQL_EDITOR_LINENUMBER_ARG} = '--line ' \]

The default is + on Unix systems (corresponding to the default editor vi, and useful for many other common editors); but there is no default on Windows systems.

Alternative location for the command history file. Tilde (~) expansion is performed.

If a query's results do not fit on the screen, they are piped through this command. Typical values are more or less. Use of the pager can be disabled by setting PSQL_PAGER or PAGER to an empty string, or by adjusting the pager-related options of the \pset command. These variables are examined in the order listed; the first that is set is used. If none of them is set, the default is to use more on most platforms, but less on Cygwin.

Alternative location of the user's .psqlrc file. Tilde (~) expansion is performed.

Command executed by the \! command.

Directory for storing temporary files. The default is /tmp.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Files

\[ \text{psqlrc and ~/.psqlrc} \]

Unless it is passed an --X option, psql attempts to read and execute commands from the system-wide startup file (psqlrc) and then the user's personal startup file (~/.psqlrc), after connecting to the database but before accepting normal commands. These files can be used to set up the client and/or the server to taste, typically with \set and SET commands.

The system-wide startup file is named psqlrc and is sought in the installation's "system configuration" directory, which is most reliably identified by running pg_config --sysconfigdir.
By default this directory will be ../etc/ relative to the directory containing the PostgreSQL executables. The name of this directory can be set explicitly via the PGSYSCONFDIR environment variable.

The user's personal startup file is named .psqlrc and is sought in the invoking user's home directory. On Windows, which lacks such a concept, the personal startup file is named %APPDATA%\postgresql\psqlrc.conf. The location of the user's startup file can be set explicitly via the PSQLRC environment variable.

Both the system-wide startup file and the user's personal startup file can be made psql-version-specific by appending a dash and the PostgreSQL major or minor release number to the file name, for example ~/.psqlrc-9.2 or ~/.psqlrc-9.2.5. The most specific version-matching file will be read in preference to a non-version-specific file.

.psql_history

The command-line history is stored in the file ~/.psql_history, or %APPDATA%\postgresql\psql_history on Windows.

The location of the history file can be set explicitly via the HISTFILE psql variable or the PSQL_HISTORY environment variable.

Notes

• psql works best with servers of the same or an older major version. Backslash commands are particularly likely to fail if the server is of a newer version than psql itself. However, backslash commands of the \d family should work with servers of versions back to 7.4, though not necessarily with servers newer than psql itself. The general functionality of running SQL commands and displaying query results should also work with servers of a newer major version, but this cannot be guaranteed in all cases.

If you want to use psql to connect to several servers of different major versions, it is recommended that you use the newest version of psql. Alternatively, you can keep around a copy of psql from each major version and be sure to use the version that matches the respective server. But in practice, this additional complication should not be necessary.

• Before PostgreSQL 9.6, the -c option implied -X (-no-psqlrc); this is no longer the case.

• Before PostgreSQL 8.4, psql allowed the first argument of a single-letter backslash command to start directly after the command, without intervening whitespace. Now, some whitespace is required.

Notes for Windows Users

psql is built as a “console application”. Since the Windows console windows use a different encoding than the rest of the system, you must take special care when using 8-bit characters within psql. If psql detects a problematic console code page, it will warn you at startup. To change the console code page, two things are necessary:

• Set the code page by entering cmd.exe /c chcp 1252. (1252 is a code page that is appropriate for German; replace it with your value.) If you are using Cygwin, you can put this command in / etc/profile.

• Set the console font to Lucida Console, because the raster font does not work with the ANSI code page.

Examples

The first example shows how to spread a command over several lines of input. Notice the changing prompt:
testdb=> CREATE TABLE my_table {
    first integer not null default 0,
    second text
};
CREATE TABLE

Now look at the table definition again:

testdb=> \d my_table
Table "public.my_table"
<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Collation</th>
<th>Nullable</th>
<th>Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>first</td>
<td>integer</td>
<td></td>
<td>not null</td>
<td>0</td>
</tr>
<tr>
<td>second</td>
<td>text</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Now we change the prompt to something more interesting:

testdb=> \set PROMPT1 '%n@%m %~%R%# '
peter@localhost testdb=>

Let's assume you have filled the table with data and want to take a look at it:

peter@localhost testdb=> SELECT * FROM my_table;
first second
------- --------
 1   one
 2   two
 3   three
 4   four
(4 rows)

You can display tables in different ways by using the \pset command:

peter@localhost testdb=> \pset border 2
Border style is 2.
peter@localhost testdb=> SELECT * FROM my_table;
+-------+--------+
| first | second |
+-------+--------+
| 1     | one    |
| 2     | two    |
| 3     | three  |
| 4     | four   |
+-------+--------+
(4 rows)

peter@localhost testdb=> \pset border 0
Border style is 0.
peter@localhost testdb=> SELECT * FROM my_table;
first second
------- -------
 1   one
 2   two
three
four
(4 rows)

peter@localhost testdb=> \pset border 1
Border style is 1.
peter@localhost testdb=> \pset format unaligned
Output format is unaligned.
peter@localhost testdb=> \pset fieldsep "",
Field separator is ",".
peter@localhost testdb=> \pset tuples_only
Showing only tuples.
peter@localhost testdb=> SELECT second, first FROM my_table;
one,1
two,2
three,3
four,4

Alternatively, use the short commands:

peter@localhost testdb=> \a \t \x
Output format is aligned.
Tuples only is off.
Expanded display is on.
peter@localhost testdb=> SELECT * FROM my_table;

When suitable, query results can be shown in a crosstab representation with the \crosstabview command:

testdb=> SELECT first, second, first > 2 AS gt2 FROM my_table;

<table>
<thead>
<tr>
<th>first</th>
<th>second</th>
<th>gt2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
<td>f</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
<td>f</td>
</tr>
<tr>
<td>3</td>
<td>three</td>
<td>t</td>
</tr>
<tr>
<td>4</td>
<td>four</td>
<td>t</td>
</tr>
</tbody>
</table>

(4 rows)

testdb=> \crosstabview first second

<table>
<thead>
<tr>
<th>first</th>
<th>one</th>
<th>two</th>
<th>three</th>
<th>four</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>f</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>f</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>t</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>t</td>
</tr>
</tbody>
</table>
This second example shows a multiplication table with rows sorted in reverse numerical order and columns with an independent, ascending numerical order.

```sql
SELECT t1.first as "A", t2.first+100 AS "B",
       t1.first*(t2.first+100) as "AxB",
       row_number() over(order by t2.first) AS ord
FROM my_table t1 CROSS JOIN my_table t2 ORDER BY 1 DESC
\crosstabview "A" "B" "AxB" ord
```

<table>
<thead>
<tr>
<th></th>
<th>101</th>
<th>102</th>
<th>103</th>
<th>104</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>404</td>
<td>408</td>
<td>412</td>
<td>416</td>
</tr>
<tr>
<td>3</td>
<td>303</td>
<td>306</td>
<td>309</td>
<td>312</td>
</tr>
<tr>
<td>2</td>
<td>202</td>
<td>204</td>
<td>206</td>
<td>208</td>
</tr>
<tr>
<td>1</td>
<td>101</td>
<td>102</td>
<td>103</td>
<td>104</td>
</tr>
</tbody>
</table>

(4 rows)
reindexdb

reindexdb — reindex a PostgreSQL database

Synopsis

reindexdb [connection-option...] [option...] [--schema|--schema=-S schema]... [--table|--table=-t table]... [--index|--index=-i index]... [dbname]

reindexdb [connection-option...] [option...|--all|-a]

reindexdb [connection-option...] [option...|--system|--system=-s [dbname]

Description

reindexdb is a utility for rebuilding indexes in a PostgreSQL database.

reindexdb is a wrapper around the SQL command REINDEX. There is no effective difference between reindexing databases via this utility and via other methods for accessing the server.

Options

reindexdb accepts the following command-line arguments:

-a
--all

Reindex all databases.

[-d] dbname
[--dbname=] dbname

Specifies the name of the database to be reindexed. If this is not specified and -a (or --all) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

-e
--echo

Echo the commands that reindexdb generates and sends to the server.

-i index
--index=index

Recreate index only. Multiple indexes can be recreated by writing multiple -i switches.

-q
--quiet

Do not display progress messages.

-s
--system

Reindex database's system catalogs.
reindexdb

-S schema
--schema=schema

Reindex schema only. Multiple schemas can be reindexed by writing multiple -S switches.

-t table
--table=table

Reindex table only. Multiple tables can be reindexed by writing multiple -t switches.

-v
--verbose

Print detailed information during processing.

-V
--version

Print the reindexdb version and exit.

-?
--help

Show help about reindexdb command line arguments, and exit.

reindexdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force reindexdb to prompt for a password before connecting to a database.

This option is never essential, since reindexdb will automatically prompt for a password if the server demands password authentication. However, reindexdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -w to avoid the extra connection attempt.
--maintenance-db=dbname

Specifies the name of the database to connect to discover what other databases should be reindexed. If not specified, the *postgres* database will be used, and if that does not exist, *template1* will be used.

**Environment**

*PGDATABASE*  
*PGHOST*  
*PGPORT*  
*PGUSER*

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

**Diagnostics**

In case of difficulty, see REINDEX and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

**Notes**

reindexdb might need to connect several times to the PostgreSQL server, asking for a password each time. It is convenient to have a `~/.pgpass` file in such cases. See Section 34.15 for more information.

**Examples**

To reindex the database *test*:

```
$ reindexdb test
```

To reindex the table *foo* and the index *bar* in a database named *abcd*:

```
$ reindexdb --table=foo --index=bar abcd
```

**See Also**

REINDEX
vacuumdb

vacuumdb — garbage-collect and analyze a PostgreSQL database

Synopsis

vacuumdb [connection-option...][option...][-table|-t table[(column[,...])]]...
[dbname]
vacuumdb [connection-option...][option...][--all|-a]

Description

vacuumdb is a utility for cleaning a PostgreSQL database. vacuumdb will also generate internal statistics used by the PostgreSQL query optimizer.

vacuumdb is a wrapper around the SQL command VACUUM. There is no effective difference between vacuuming and analyzing databases via this utility and via other methods for accessing the server.

Options

vacuumdb accepts the following command-line arguments:

- -a
  --all
  Vacuum all databases.

[-d] dbname
[--dbname=] dbname

  Specifies the name of the database to be cleaned or analyzed. If this is not specified and -a (or --all) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

- -e
  --echo

  Echo the commands that vacuumdb generates and sends to the server.

- f
  --full

  Perform “full” vacuuming.

- F
  --freeze

  Aggressively “freeze” tuples.

- j njobs
  --jobs=njobs

  Execute the vacuum or analyze commands in parallel by running njobs commands simultaneously. This option reduces the time of the processing but it also increases the load on the database server.
vacuumdb will open $n$ jobs connections to the database, so make sure your max_connections setting is high enough to accommodate all connections.

Note that using this mode together with the -f (FULL) option might cause deadlock failures if certain system catalogs are processed in parallel.

-q
--quiet

Do not display progress messages.

-t table [ (column [, ...]) ]
--table=table [ (column [, ...]) ]

Clean or analyze table only. Column names can be specified only in conjunction with the --analyze or --analyze-only options. Multiple tables can be vacuumed by writing multiple -t switches.

Tip

If you specify columns, you probably have to escape the parentheses from the shell. (See examples below.)

-v
--verbose

Print detailed information during processing.

-V
--version

Print the vacuumdb version and exit.

-z
--analyze

Also calculate statistics for use by the optimizer.

-Z
--analyze-only

Only calculate statistics for use by the optimizer (no vacuum).

--analyze-in-stages

Only calculate statistics for use by the optimizer (no vacuum), like --analyze-only. Run several (currently three) stages of analyze with different configuration settings, to produce usable statistics faster.

This option is useful to analyze a database that was newly populated from a restored dump or by pg_upgrade. This option will try to create some statistics as fast as possible, to make the database usable, and then produce full statistics in the subsequent stages.

-?
--help

Show help about vacuumdb command line arguments, and exit.
vacuumdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force vacuumdb to prompt for a password before connecting to a database.

This option is never essential, since vacuumdb will automatically prompt for a password if the server demands password authentication. However, vacuumdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to discover what other databases should be vacu-umed. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.

Environment

PGDATABASE
PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

Diagnostics

In case of difficulty, see VACUUM and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.
vacuumdb might need to connect several times to the PostgreSQL server, asking for a password each time. It is convenient to have a ~/.pgpass file in such cases. See Section 34.15 for more information.

Examples

To clean the database test:

$ vacuumdb test

To clean and analyze for the optimizer a database named bigdb:

$ vacuumdb --analyze bigdb

To clean a single table foo in a database named xyzzy, and analyze a single column bar of the table for the optimizer:

$ vacuumdb --analyze --verbose --table='foo(bar)' xyzzy

See Also

VACUUM
PostgreSQL Server Applications

This part contains reference information for PostgreSQL server applications and support utilities. These commands can only be run usefully on the host where the database server resides. Other utility programs are listed in PostgreSQL Client Applications.

Table of Contents

initdb ................................................................. 1945
pg_archivecleanup ................................................. 1949
pg_controldata .................................................... 1951
pg_ctl ............................................................... 1952
pg_resetwal ......................................................... 1958
pg_rollback ......................................................... 1961
pg_test_fsync .................................................... 1964
pg_test_timing ...................................................... 1965
pg_upgrade .......................................................... 1969
pg_verify_checksums ................................................ 1977
pg_waldump ........................................................ 1978
postgres ............................................................. 1980
postmaster .......................................................... 1987
initdb

initdb — create a new PostgreSQL database cluster

Synopsis

initdb [option...][ --pgdata | -D ] directory

Description

initdb creates a new PostgreSQL database cluster. A database cluster is a collection of databases that are managed by a single server instance.

Creating a database cluster consists of creating the directories in which the database data will live, generating the shared catalog tables (tables that belong to the whole cluster rather than to any particular database), and creating the template1 and postgres databases. When you later create a new database, everything in the template1 database is copied. (Therefore, anything installed in template1 is automatically copied into each database created later.) The postgres database is a default database meant for use by users, utilities and third party applications.

Although initdb will attempt to create the specified data directory, it might not have permission if the parent directory of the desired data directory is root-owned. To initialize in such a setup, create an empty data directory as root, then use chown to assign ownership of that directory to the database user account, then su to become the database user to run initdb.

initdb must be run as the user that will own the server process, because the server needs to have access to the files and directories that initdb creates. Since the server cannot be run as root, you must not run initdb as root either. (It will in fact refuse to do so.)

For security reasons the new cluster created by initdb will only be accessible by the cluster owner by default. The --allow-group-access option allows any user in the same group as the cluster owner to read files in the cluster. This is useful for performing backups as a non-privileged user.

initdb initializes the database cluster's default locale and character set encoding. The character set encoding, collation order (LC_COLLATE) and character set classes (LC_CTYPE, e.g. upper, lower, digit) can be set separately for a database when it is created. initdb determines those settings for the template1 database, which will serve as the default for all other databases.

To alter the default collation order or character set classes, use the --lc-collate and --lc-ctype options. Collation orders other than C or POSIX also have a performance penalty. For these reasons it is important to choose the right locale when running initdb.

The remaining locale categories can be changed later when the server is started. You can also use --locale to set the default for all locale categories, including collation order and character set classes. All server locale values (lc_*) can be displayed via SHOW ALL. More details can be found in Section 23.1.

To alter the default encoding, use the --encoding. More details can be found in Section 23.3.

Options

-A authmethod
--auth=authmethod

This option specifies the default authentication method for local users used in pg_hba.conf (host and local lines). initdb will prepopulate pg_hba.conf entries using the specified authentication method for non-replication as well as replication connections.
Do not use `trust` unless you trust all local users on your system. `trust` is the default for ease of installation.

```bash
--auth-host=authmethod
```

This option specifies the authentication method for local users via TCP/IP connections used in `pg_hba.conf` (host lines).

```bash
--auth-local=authmethod
```

This option specifies the authentication method for local users via Unix-domain socket connections used in `pg_hba.conf` (local lines).

```bash
-D directory
--pgdata=directory
```

This option specifies the directory where the database cluster should be stored. This is the only information required by `initdb`, but you can avoid writing it by setting the `PGDATA` environment variable, which can be convenient since the database server (`postgres`) can find the database directory later by the same variable.

```bash
-E encoding
--encoding=encoding
```

Selects the encoding of the template database. This will also be the default encoding of any database you create later, unless you override it there. The default is derived from the locale, or `SQL_ASCII` if that does not work. The character sets supported by the PostgreSQL server are described in Section 23.3.1.

```bash
-g
--allow-group-access
```

Allows users in the same group as the cluster owner to read all cluster files created by `initdb`.

```bash
-k
--data-checksums
```

Use checksums on data pages to help detect corruption by the I/O system that would otherwise be silent. Enabling checksums may incur a noticeable performance penalty. This option can only be set during initialization, and cannot be changed later. If set, checksums are calculated for all objects, in all databases.

```bash
--locale=locale
```

Sets the default locale for the database cluster. If this option is not specified, the locale is inherited from the environment that `initdb` runs in. Locale support is described in Section 23.1.

```bash
--lc-collate=locale
--lc-ctype=locale
--lc-messages=locale
--lc-monetary=locale
--lc-numeric=locale
--lc-time=locale
```

Like `--locale`, but only sets the locale in the specified category.

```bash
--no-locale
```

Equivalent to `--locale=C`.

1946
By default, initdb will wait for all files to be written safely to disk. This option causes initdb to return without waiting, which is faster, but means that a subsequent operating system crash can leave the data directory corrupt. Generally, this option is useful for testing, but should not be used when creating a production installation.

--pwfile=filename

Makes initdb read the database superuser's password from a file. The first line of the file is taken as the password.

-S
--sync-only

Safely write all database files to disk and exit. This does not perform any of the normal initdb operations.

-T config
--text-search-config=config

Sets the default text search configuration. See default_text_search_config for further information.

-U username
--username=username

Selects the user name of the database superuser. This defaults to the name of the effective user running initdb. It is really not important what the superuser's name is, but one might choose to keep the customary name postgres, even if the operating system user's name is different.

-W
--pwprompt

Makes initdb prompt for a password to give the database superuser. If you don't plan on using password authentication, this is not important. Otherwise you won't be able to use password authentication until you have a password set up.

-X directory
--waldir=directory

This option specifies the directory where the write-ahead log should be stored.

--wal-segsize=size

Set the WAL segment size, in megabytes. This is the size of each individual file in the WAL log. The default size is 16 megabytes. The value must be a power of 2 between 1 and 1024 (megabytes). This option can only be set during initialization, and cannot be changed later.

It may be useful to adjust this size to control the granularity of WAL log shipping or archiving. Also, in databases with a high volume of WAL, the sheer number of WAL files per directory can become a performance and management problem. Increasing the WAL file size will reduce the number of WAL files.

Other, less commonly used, options are also available:

-d
--debug

Print debugging output from the bootstrap backend and a few other messages of lesser interest for the general public. The bootstrap backend is the program initdb uses to create the catalog tables. This option generates a tremendous amount of extremely boring output.
-L directory

Specifies where \texttt{initdb} should find its input files to initialize the database cluster. This is normally not necessary. You will be told if you need to specify their location explicitly.

-n
--no-clean

By default, when \texttt{initdb} determines that an error prevented it from completely creating the database cluster, it removes any files it might have created before discovering that it cannot finish the job. This option inhibits tidying-up and is thus useful for debugging.

Other options:

-V
--version

Print the \texttt{initdb} version and exit.

-?
--help

Show help about \texttt{initdb} command line arguments, and exit.

**Environment**

\texttt{PGDATA}

Specifies the directory where the database cluster is to be stored; can be overridden using the \texttt{-D} option.

\texttt{TZ}

Specifies the default time zone of the created database cluster. The value should be a full time zone name (see Section 8.5.3).

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

**Notes**

\texttt{initdb} can also be invoked via \texttt{pg_ctl initdb}.

**See Also**

\texttt{pg_ctl}, \texttt{postgres}
pg_archivecleanup

pg_archivecleanup — clean up PostgreSQL WAL archive files

Synopsis

pg_archivecleanup [option...] archivelocation oldestkeptwalfile

Description

pg_archivecleanup is designed to be used as an archive_cleanup_command to clean up WAL file archives when running as a standby server (see Section 26.2). pg_archivecleanup can also be used as a standalone program to clean WAL file archives.

To configure a standby server to use pg_archivecleanup, put this into its recovery.conf configuration file:

archive_cleanup_command = 'pg_archivecleanup archivelocation %r'

where archivelocation is the directory from which WAL segment files should be removed.

When used within archive_cleanup_command, all WAL files logically preceding the value of the %r argument will be removed from archivelocation. This minimizes the number of files that need to be retained, while preserving crash-restart capability. Use of this parameter is appropriate if the archivelocation is a transient staging area for this particular standby server, but not when the archivelocation is intended as a long-term WAL archive area, or when multiple standby servers are recovering from the same archive location.

When used as a standalone program all WAL files logically preceding the oldestkeptwalfile will be removed from archivelocation. In this mode, if you specify a .partial or .backup file name, then only the file prefix will be used as the oldestkeptwalfile. This treatment of .backup file name allows you to remove all WAL files archived prior to a specific base backup without error. For example, the following example will remove all files older than WAL file name 00000010000003700000010:

pg_archivecleanup -d archive
000000010000003700000010.00000020.backup

pg_archivecleanup: keep WAL file
"archive/000000010000003700000010" and later
pg_archivecleanup: removing file
"archive/00000001000000370000000F"
pg_archivecleanup: removing file
"archive/00000001000000370000000E"

pg_archivecleanup assumes that archivelocation is a directory readable and writable by the server-owning user.

Options

pg_archivecleanup accepts the following command-line arguments:

-d

Print lots of debug logging output on stderr.
pg_archivecleanup

-\n
Print the names of the files that would have been removed on stdout (performs a dry run).

-V
--version

Print the pg_archivecleanup version and exit.

-x extension

Provide an extension that will be stripped from all file names before deciding if they should be deleted. This is typically useful for cleaning up archives that have been compressed during storage, and therefore have had an extension added by the compression program. For example: -x .gz.

-?
--help

Show help about pg_archivecleanup command line arguments, and exit.

Notes

pg_archivecleanup is designed to work with PostgreSQL 8.0 and later when used as a standalone utility, or with PostgreSQL 9.0 and later when used as an archive cleanup command.

pg_archivecleanup is written in C and has an easy-to-modify source code, with specifically designated sections to modify for your own needs.

Examples

On Linux or Unix systems, you might use:

```
archive_cleanup_command = 'pg_archivecleanup -d /mnt/standby/archive %r 2>>cleanup.log'
```

where the archive directory is physically located on the standby server, so that the archive_command is accessing it across NFS, but the files are local to the standby. This will:

- produce debugging output in cleanup.log
- remove no-longer-needed files from the archive directory

See Also

pg_standby
pg_controldata

pg_controldata — display control information of a PostgreSQL database cluster

Synopsis

pg_controldata [option] [ --pgdata | -D ] datadir

Description

pg_controldata prints information initialized during initdb, such as the catalog version. It also shows information about write-ahead logging and checkpoint processing. This information is cluster-wide, and not specific to any one database.

This utility can only be run by the user who initialized the cluster because it requires read access to the data directory. You can specify the data directory on the command line, or use the environment variable PGDATA. This utility supports the options -V and --version, which print the pg_controldata version and exit. It also supports options -? and --help, which output the supported arguments.

Environment

PGDATA

Default data directory location
pg_ctl

pg_ctl — initialize, start, stop, or control a PostgreSQL server

Synopsis

pg_ctl init [db] [-D datadir] [-s] [-o initdb-options]
pg_ctl reload [-D datadir] [-s]
pg_ctl status [-D datadir]
pg_ctl promote [-D datadir] [-W] [-t seconds] [-s]
pg_ctl kill signal_name process_id

On Microsoft Windows, also:

pg_ctl unregister [-N servicename]

Description

pg_ctl is a utility for initializing a PostgreSQL database cluster, starting, stopping, or restarting the PostgreSQL database server (postgres), or displaying the status of a running server. Although the server can be started manually, pg_ctl encapsulates tasks such as redirecting log output and properly detaching from the terminal and process group. It also provides convenient options for controlled shutdown.

The init or initdb mode creates a new PostgreSQL database cluster, that is, a collection of databases that will be managed by a single server instance. This mode invokes the initdb command. See initdb for details.

start mode launches a new server. The server is started in the background, and its standard input is attached to /dev/null (or nul on Windows). On Unix-like systems, by default, the server's standard output and standard error are sent to pg_ctl's standard output (not standard error). The standard output of pg_ctl should then be redirected to a file or piped to another process such as a log rotating program like rotatelogs; otherwise postgres will write its output to the controlling terminal (from the background) and will not leave the shell's process group. On Windows, by default the server's standard output and standard error are sent to the terminal. These default behaviors can be changed by using -l to append the server's output to a log file. Use of either -l or output redirection is recommended.

stop mode shuts down the server that is running in the specified data directory. Three different shutdown methods can be selected with the -m option. “Smart” mode waits for all active clients to
disconnect and any online backup to finish. If the server is in hot standby, recovery and streaming replication will be terminated once all clients have disconnected. “Fast” mode (the default) does not wait for clients to disconnect and will terminate an online backup in progress. All active transactions are rolled back and clients are forcibly disconnected, then the server is shut down. “Immediate” mode will abort all server processes immediately, without a clean shutdown. This choice will lead to a crash-recovery cycle during the next server start.

restart mode effectively executes a stop followed by a start. This allows changing the postgres command-line options, or changing configuration-file options that cannot be changed without restarting the server. If relative paths were used on the command line during server start, restart might fail unless pg_ctl is executed in the same current directory as it was during server start.

reload mode simply sends the postgres server process a SIGHUP signal, causing it to reread its configuration files (postgresql.conf, pg_hba.conf, etc.). This allows changing configuration-file options that do not require a full server restart to take effect.

status mode checks whether a server is running in the specified data directory. If it is, the server’s PID and the command line options that were used to invoke it are displayed. If the server is not running, pg_ctl returns an exit status of 3. If an accessible data directory is not specified, pg_ctl returns an exit status of 4.

promote mode commands the standby server that is running in the specified data directory to end standby mode and begin read-write operations.

kill mode sends a signal to a specified process. This is primarily valuable on Microsoft Windows which does not have a built-in kill command. Use --help to see a list of supported signal names.

register mode registers the PostgreSQL server as a system service on Microsoft Windows. The -S option allows selection of service start type, either “auto” (start service automatically on system startup) or “demand” (start service on demand).

unregister mode unregisters a system service on Microsoft Windows. This undoes the effects of the register command.

Options

-c
--core-files

Attempt to allow server crashes to produce core files, on platforms where this is possible, by lifting any soft resource limit placed on core files. This is useful in debugging or diagnosing problems by allowing a stack trace to be obtained from a failed server process.

-D datadir
--pgdata=datadir

Specifies the file system location of the database configuration files. If this option is omitted, the environment variable PGDATA is used.

-l filename
--log=filename

Append the server log output to filename. If the file does not exist, it is created. The umask is set to 077, so access to the log file is disallowed to other users by default.

-m mode
--mode=mode

Specifies the shutdown mode. mode can be smart, fast, or immediate, or the first letter of one of these three. If this option is omitted, fast is the default.
-o options
--options=options

Specifies options to be passed directly to the postgres command. -o can be specified multiple times, with all the given options being passed through.

The options should usually be surrounded by single or double quotes to ensure that they are passed through as a group.

-o initdb-options
--options=initdb-options

Specifies options to be passed directly to the initdb command. -o can be specified multiple times, with all the given options being passed through.

The initdb-options should usually be surrounded by single or double quotes to ensure that they are passed through as a group.

-p path

Specifies the location of the postgres executable. By default the postgres executable is taken from the same directory as pg_ctl, or failing that, the hard-wired installation directory. It is not necessary to use this option unless you are doing something unusual and get errors that the postgres executable was not found.

In init mode, this option analogously specifies the location of the initdb executable.

-s
--silent

Print only errors, no informational messages.

-t seconds
--timeout=seconds

Specifies the maximum number of seconds to wait when waiting for an operation to complete (see option -w). Defaults to the value of the PGCTLS_TIMEOUT environment variable or, if not set, to 60 seconds.

-V
--version

Print the pg_ctl version and exit.

-w
--wait

Wait for the operation to complete. This is supported for the modes start, stop, restart, promote, and register, and is the default for those modes.

When waiting, pg_ctl repeatedly checks the server’s PID file, sleeping for a short amount of time between checks. Startup is considered complete when the PID file indicates that the server is ready to accept connections. Shutdown is considered complete when the server removes the PID file. pg_ctl returns an exit code based on the success of the startup or shutdown.

If the operation does not complete within the timeout (see option -t), then pg_ctl exits with a nonzero exit status. But note that the operation might continue in the background and eventually succeed.

-W
--no-wait

Do not wait for the operation to complete. This is the opposite of the option -w.
If waiting is disabled, the requested action is triggered, but there is no feedback about its success. In that case, the server log file or an external monitoring system would have to be used to check the progress and success of the operation.

In prior releases of PostgreSQL, this was the default except for the stop mode.

-?  help
Show help about pg_ctl command line arguments, and exit.

If an option is specified that is valid, but not relevant to the selected operating mode, pg_ctl ignores it.

**Options for Windows**

- e source
Name of the event source for pg_ctl to use for logging to the event log when running as a Windows service. The default is PostgreSQL. Note that this only controls messages sent from pg_ctl itself; once started, the server will use the event source specified by its event_source parameter. Should the server fail very early in startup, before that parameter has been set, it might also log using the default event source name PostgreSQL.

-N servicename
Name of the system service to register. This name will be used as both the service name and the display name. The default is PostgreSQL.

-P password
Password for the user to run the service as.

-S start-type
Start type of the system service. start-type can be auto, or demand, or the first letter of one of these two. If this option is omitted, auto is the default.

-U username
User name for the user to run the service as. For domain users, use the format DOMAIN\username.

**Environment**

PGCTLTIMEOUT
Default limit on the number of seconds to wait when waiting for startup or shutdown to complete. If not set, the default is 60 seconds.

PGDATA
Default data directory location.

Most pg_ctl modes require knowing the data directory location; therefore, the -D option is required unless PGDATA is set.

pg_ctl, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 34.14).

For additional variables that affect the server, see postgres.
Files

postmaster.pid

pg_ctl examines this file in the data directory to determine whether the server is currently running.

postmaster.opts

If this file exists in the data directory, pg_ctl (in restart mode) will pass the contents of the file as options to postgres, unless overridden by the -o option. The contents of this file are also displayed in status mode.

Examples

Starting the Server

To start the server, waiting until the server is accepting connections:

$ pg_ctl start

To start the server using port 5433, and running without fsync, use:

$ pg_ctl -o "-F -p 5433" start

Stopping the Server

To stop the server, use:

$ pg_ctl stop

The -m option allows control over how the server shuts down:

$ pg_ctl stop -m smart

Restarting the Server

Restarting the server is almost equivalent to stopping the server and starting it again, except that by default, pg_ctl saves and reuses the command line options that were passed to the previously-running instance. To restart the server using the same options as before, use:

$ pg_ctl restart

But if -o is specified, that replaces any previous options. To restart using port 5433, disabling fsync upon restart:

$ pg_ctl -o "-F -p 5433" restart

Showing the Server Status

Here is sample status output from pg_ctl:

$ pg_ctl status
pg_ctl

pg_ctl: server is running (PID: 13718)
/usr/local/pgsql/bin/postgres "-D" "/usr/local/pgsql/data" "-p" 
"5433" "-B" "128"

The second line is the command that would be invoked in restart mode.

See Also

initdb, postgres
pg_resetwal

pg_resetwal — reset the write-ahead log and other control information of a PostgreSQL database cluster

Synopsis

pg_resetwal[--force|-f][--dry-run|-n][option...][--pgdata|-D] datadir

Description

pg_resetwal clears the write-ahead log (WAL) and optionally resets some other control information stored in the pg_control file. This function is sometimes needed if these files have become corrupted. It should be used only as a last resort, when the server will not start due to such corruption.

After running this command, it should be possible to start the server, but bear in mind that the database might contain inconsistent data due to partially-committed transactions. You should immediately dump your data, run initdb, and reload. After reload, check for inconsistencies and repair as needed.

This utility can only be run by the user who installed the server, because it requires read/write access to the data directory. For safety reasons, you must specify the data directory on the command line. pg_resetwal does not use the environment variable PGDATA.

If pg_resetwal complains that it cannot determine valid data for pg_control, you can force it to proceed anyway by specifying the -f (force) option. In this case plausible values will be substituted for the missing data. Most of the fields can be expected to match, but manual assistance might be needed for the next OID, next transaction ID and epoch, next multitransaction ID and offset, and WAL starting location fields. These fields can be set using the options discussed below. If you are not able to determine correct values for all these fields, -f can still be used, but the recovered database must be treated with even more suspicion than usual: an immediate dump and reload is imperative. Do not execute any data-modifying operations in the database before you dump, as any such action is likely to make the corruption worse.

Options

- -f
--force

Force pg_resetwal to proceed even if it cannot determine valid data for pg_control, as explained above.

- -n
--dry-run

The -n/--dry-run option instructs pg_resetwal to print the values reconstructed from pg_control and values about to be changed, and then exit without modifying anything. This is mainly a debugging tool, but can be useful as a sanity check before allowing pg_resetwal to proceed for real.

- -V
--version

Display version information, then exit.

- -?
--help

Show help, then exit.
The following options are only needed when `pg_resetwal` is unable to determine appropriate values by reading `pg_control`. Safe values can be determined as described below. For values that take numeric arguments, hexadecimal values can be specified by using the prefix `0x`.

- **-c xid,xid**
  --commit-timestamp-ids=xid,xid

  Manually set the oldest and newest transaction IDs for which the commit time can be retrieved.

  A safe value for the oldest transaction ID for which the commit time can be retrieved (first part) can be determined by looking for the numerically smallest file name in the directory `pg_commit_ts` under the data directory. Conversely, a safe value for the newest transaction ID for which the commit time can be retrieved (second part) can be determined by looking for the numerically greatest file name in the same directory. The file names are in hexadecimal.

- **-e xid_epoch**
  --epoch=xid_epoch

  Manually set the next transaction ID's epoch.

  The transaction ID epoch is not actually stored anywhere in the database except in the field that is set by `pg_resetwal`, so any value will work so far as the database itself is concerned. You might need to adjust this value to ensure that replication systems such as Slony-I and Skytools work correctly — if so, an appropriate value should be obtainable from the state of the downstream replicated database.

- **-l walfile**
  --next-wal-file=walfile

  Manually set the WAL starting location by specifying the name of the next WAL segment file.

  The name of next WAL segment file should be larger than any WAL segment file name currently existing in the directory `pg_wal` under the data directory. These names are also in hexadecimal and have three parts. The first part is the “timeline ID” and should usually be kept the same. For example, if `00000001000000320000004A` is the largest entry in `pg_wal`, use `-l 00000001000000320000004B` or higher.

  Note that when using nondefault WAL segment sizes, the numbers in the WAL file names are different from the LSNs that are reported by system functions and system views. This option takes a WAL file name, not an LSN.

  **Note**

  `pg_resetwal` itself looks at the files in `pg_wal` and chooses a default `-l` setting beyond the last existing file name. Therefore, manual adjustment of `-l` should only be needed if you are aware of WAL segment files that are not currently present in `pg_wal`, such as entries in an offline archive; or if the contents of `pg_wal` have been lost entirely.

- **-m mxid,mxid**
  --multixact-ids=mxid,mxid

  Manually set the next and oldest multitransaction ID.

  A safe value for the next multitransaction ID (first part) can be determined by looking for the numerically largest file name in the directory `pg_multixact/offsets` under the data directory, adding one, and then multiplying by 65536 (0x10000). Conversely, a safe value for the oldest multitransaction ID (second part of `-m`) can be determined by looking for the numerically smallest
file name in the same directory and multiplying by 65536. The file names are in hexadecimal, so the easiest way to do this is to specify the option value in hexadecimal and append four zeroes.

```
-o oid
--next-oid=oid
```

Manually set the next OID.

There is no comparably easy way to determine a next OID that's beyond the largest one in the database, but fortunately it is not critical to get the next-OID setting right.

```
-O mxoff
--multixact-offset=mxoff
```

Manually set the next multitransaction offset.

A safe value can be determined by looking for the numerically largest file name in the directory `pg_multixact/members` under the data directory, adding one, and then multiplying by 52352 (0xCC80). The file names are in hexadecimal. There is no simple recipe such as the ones for other options of appending zeroes.

```
--wal-segsize=wal_segment_size
```

Set the new WAL segment size, in megabytes. The value must be set to a power of 2 between 1 and 1024 (megabytes). See the same option of initdb for more information.

```
Note

While `pg_resetwal` will set the WAL starting address beyond the latest existing WAL segment file, some segment size changes can cause previous WAL file names to be reused. It is recommended to use `-l` together with this option to manually set the WAL starting address if WAL file name overlap will cause problems with your archiving strategy.
```

```
-x xid
--next-transaction-id=xid
```

Manually set the next transaction ID.

A safe value can be determined by looking for the numerically largest file name in the directory `pg_xact` under the data directory, adding one, and then multiplying by 1048576 (0x100000). Note that the file names are in hexadecimal. It is usually easiest to specify the option value in hexadecimal too. For example, if `0011` is the largest entry in `pg_xact`, `--next-transaction-id=0x1200000` will work (five trailing zeroes provide the proper multiplier).

```
Notes

This command must not be used when the server is running. `pg_resetwal` will refuse to start up if it finds a server lock file in the data directory. If the server crashed then a lock file might have been left behind; in that case you can remove the lock file to allow `pg_resetwal` to run. But before you do so, make doubly certain that there is no server process still alive.

`pg_resetwal` works only with servers of the same major version.

```
See Also

`pg_controldata`
```
**pg_rewind**

*pg_rewind* — synchronize a PostgreSQL data directory with another data directory that was forked from it

**Synopsis**

```bash
go_rewind [option...] {-D | --target-pgdata} directory { --source-pgdata=directory | --source-server=connstr }
```

**Description**

*pg_rewind* is a tool for synchronizing a PostgreSQL cluster with another copy of the same cluster, after the clusters' timelines have diverged. A typical scenario is to bring an old master server back online after failover as a standby that follows the new master.

The result is equivalent to replacing the target data directory with the source one. Only changed blocks from relation files are copied; all other files are copied in full, including configuration files. The advantage of *pg_rewind* over taking a new base backup, or tools like rsync, is that *pg_rewind* does not require reading through unchanged blocks in the cluster. This makes it a lot faster when the database is large and only a small fraction of blocks differ between the clusters.

*pg_rewind* examines the timeline histories of the source and target clusters to determine the point where they diverged, and expects to find WAL in the target cluster's *pg_wal* directory reaching all the way back to the point of divergence. The point of divergence can be found either on the target timeline, the source timeline, or their common ancestor. In the typical failover scenario where the target cluster was shut down soon after the divergence, this is not a problem, but if the target cluster ran for a long time after the divergence, the old WAL files might no longer be present. In that case, they can be manually copied from the WAL archive to the *pg_wal* directory, or fetched on startup by configuring *recovery.conf*. The use of *pg_rewind* is not limited to failover, e.g. a standby server can be promoted, run some write transactions, and then rewinded to become a standby again.

When the target server is started for the first time after running *pg_rewind*, it will go into recovery mode and replay all WAL generated in the source server after the point of divergence. If some of the WAL was no longer available in the source server when *pg_rewind* was run, and therefore could not be copied by the *pg_rewind* session, it must be made available when the target server is started. This can be done by creating a *recovery.conf* file in the target data directory with a suitable *restore_command*.

*pg_rewind* requires that the target server either has the *wal_log_hints* option enabled in *postgresql.conf* or data checksums enabled when the cluster was initialized with *initdb*. Neither of these are currently on by default. *full_page_writes* must also be set to *on*, but is enabled by default.

**Warning**

If *pg_rewind* fails while processing, then the data folder of the target is likely not in a state that can be recovered. In such a case, taking a new fresh backup is recommended.

*pg_rewind* will fail immediately if it finds files it cannot write directly to. This can happen for example when the source and the target server use the same file mapping for read-only SSL keys and certificates. If such files are present on the target server it is recommended to remove them before running *pg_rewind*. After doing the rewind, some of those files may have been copied from the source, in which case it may be necessary to remove the data copied and restore back the set of links used before the rewind.
Options

pg_rewind accepts the following command-line arguments:

-D directory
--target-pgdata=directory

This option specifies the target data directory that is synchronized with the source. The target server must be shut down cleanly before running pg_rewind.

--source-pgdata=directory

Specifies the file system path to the data directory of the source server to synchronize the target with. This option requires the source server to be cleanly shut down.

--source-server=connstr

Specifies a libpq connection string to connect to the source PostgreSQL server to synchronize the target with. The connection must be a normal (non-replication) connection with superuser access. This option requires the source server to be running and not in recovery mode.

-n
--dry-run

Do everything except actually modifying the target directory.

-P
--progress

Enables progress reporting. Turning this on will deliver an approximate progress report while copying data from the source cluster.

--debug

Print verbose debugging output that is mostly useful for developers debugging pg_rewind.

-V
--version

Display version information, then exit.

-?
--help

Show help, then exit.

Environment

When --source-server option is used, pg_rewind also uses the environment variables supported by libpq (see Section 34.14).

Notes

How it works

The basic idea is to copy all file system-level changes from the source cluster to the target cluster:

1. Scan the WAL log of the target cluster, starting from the last checkpoint before the point where the source cluster's timeline history forked off from the target cluster. For each WAL record,
pg_rewind

record each data block that was touched. This yields a list of all the data blocks that were changed in the target cluster, after the source cluster forked off.

2. Copy all those changed blocks from the source cluster to the target cluster, either using direct file system access (--source-pgdata) or SQL (--source-server).

3. Copy all other files such as pg_xact and configuration files from the source cluster to the target cluster (everything except the relation files). Similarly to base backups, the contents of the directories pg_dynshmem/, pg_notify/, pg_replslot/, pg_serial/, pg_snapshots/, pg_stat_tmp/, and pg_subtrans/ are omitted from the data copied from the source cluster. Any file or directory beginning with pgsq1_tmp is omitted, as well as are backup_label, tablespace_map, pg_internal.init, postmaster.opts and postmaster.pid.

4. Apply the WAL from the source cluster, starting from the checkpoint created at failover. (Strictly speaking, pg_rewind doesn't apply the WAL, it just creates a backup label file that makes PostgreSQL start by replaying all WAL from that checkpoint forward.)
pg_test_fsync

pg_test_fsync — determine fastest wal_sync_method for PostgreSQL

Synopsis

pg_test_fsync [option...]

Description

pg_test_fsync is intended to give you a reasonable idea of what the fastest wal_sync_method is on your specific system, as well as supplying diagnostic information in the event of an identified I/O problem. However, differences shown by pg_test_fsync might not make any significant difference in real database throughput, especially since many database servers are not speed-limited by their write-ahead logs. pg_test_fsync reports average file sync operation time in microseconds for each wal_sync_method, which can also be used to inform efforts to optimize the value of commit_delay.

Options

pg_test_fsync accepts the following command-line options:

- f
  --filename

  Specifies the file name to write test data in. This file should be in the same file system that the pg_wal directory is or will be placed in. (pg_wal contains the WAL files.) The default is pg_test_fsync.out in the current directory.

- s
  --secs-per-test

  Specifies the number of seconds for each test. The more time per test, the greater the test’s accuracy, but the longer it takes to run. The default is 5 seconds, which allows the program to complete in under 2 minutes.

- V
  --version

  Print the pg_test_fsync version and exit.

- ?
  --help

  Show help about pg_test_fsync command line arguments, and exit.

See Also

postgres
pg_test_timing

pg_test_timing — measure timing overhead

Synopsis

pg_test_timing [option...]

Description

pg_test_timing is a tool to measure the timing overhead on your system and confirm that the system


time never moves backwards. Systems that are slow to collect timing data can give less accurate

EXPLAIN ANALYZE results.

Options

pg_test_timing accepts the following command-line options:

-d duration
--duration=duration

Specifies the test duration, in seconds. Longer durations give slightly better accuracy, and are more
likely to discover problems with the system clock moving backwards. The default test duration
is 3 seconds.

-V
--version

Print the pg_test_timing version and exit.

-?
--help

Show help about pg_test_timing command line arguments, and exit.

Usage

Interpreting results

Good results will show most (>90%) individual timing calls take less than one microsecond. Average
per loop overhead will be even lower, below 100 nanoseconds. This example from an Intel i7-860
system using a TSC clock source shows excellent performance:

Testing timing overhead for 3 seconds.
Per loop time including overhead: 35.96 ns
Histogram of timing durations:
< us  % of total  count
  1  96.40465   80435604
  2   3.59518   2999652
  4   0.00015     126
  8   0.00002      13
16  0.00000      2

Note that different units are used for the per loop time than the histogram. The loop can have resolution
within a few nanoseconds (ns), while the individual timing calls can only resolve down to one
microsecond (us).
Measuring executor timing overhead

When the query executor is running a statement using EXPLAIN ANALYZE, individual operations are timed as well as showing a summary. The overhead of your system can be checked by counting rows with the psql program:

```
CREATE TABLE t AS SELECT * FROM generate_series(1,100000);
\timing
SELECT COUNT(*) FROM t;
EXPLAIN ANALYZE SELECT COUNT(*) FROM t;
```

The i7-860 system measured runs the count query in 9.8 ms while the EXPLAIN ANALYZE version takes 16.6 ms, each processing just over 100,000 rows. That 6.8 ms difference means the timing overhead per row is 68 ns, about twice what pg_test_timing estimated it would be. Even that relatively small amount of overhead is making the fully timed count statement take almost 70% longer. On more substantial queries, the timing overhead would be less problematic.

Changing time sources

On some newer Linux systems, it's possible to change the clock source used to collect timing data at any time. A second example shows the slowdown possible from switching to the slower acpi_pm time source, on the same system used for the fast results above:

```
# cat /sys/devices/system/clocksource/clocksource0/available_clocksource
  tsc hpet acpi_pm
# echo acpi_pm > /sys/devices/system/clocksource/clocksource0/current_clocksource
# pg_test_timing
Per loop time including overhead: 722.92 ns
Histogram of timing durations:
< us % of total count
 1 27.84870 1155682
 2 72.05956 2990371
 4 0.07810 3241
 8 0.01357 563
16 0.00007 3
```

In this configuration, the sample EXPLAIN ANALYZE above takes 115.9 ms. That's 1061 ns of timing overhead, again a small multiple of what's measured directly by this utility. That much timing overhead means the actual query itself is only taking a tiny fraction of the accounted for time, most of it is being consumed in overhead instead. In this configuration, any EXPLAIN ANALYZE totals involving many timed operations would be inflated significantly by timing overhead.

FreeBSD also allows changing the time source on the fly, and it logs information about the timer selected during boot:

```
# dmesg | grep "Timecounter"
Timecounter "ACPI-fast" frequency 3579545 Hz quality 900
Timecounter "i8254" frequency 1193182 Hz quality 0
Timecounters tick every 10.000 msec
Timecounter "TSC" frequency 2531787134 Hz quality 800
# sysctl kern.timecounter.hardware=TSC
kern.timecounter.hardware: ACPI-fast -> TSC
```

Other systems may only allow setting the time source on boot. On older Linux systems the "clock" kernel setting is the only way to make this sort of change. And even on some more recent ones,
the only option you'll see for a clock source is "jiffies". Jiffies are the older Linux software clock implementation, which can have good resolution when it's backed by fast enough timing hardware, as in this example:

```
$ cat /sys/devices/system/clocksource/clocksource0/available_clocksource
jiffies
$ dmesg | grep time.c
time.c: Using 3.579545 MHz WALL PM GTOD PIT/TSC timer.
time.c: Detected 2400.153 MHz processor.
$ pg_test_timing
Testing timing overhead for 3 seconds.
Per timing duration including loop overhead: 97.75 ns
Histogram of timing durations:
< us   % of total      count
 1     90.23734   27694571
 2      9.75277    2993204
 4      0.00981       3010
 8      0.00007         22
16      0.00000          1
32      0.00000          1
```

### Clock hardware and timing accuracy

Collecting accurate timing information is normally done on computers using hardware clocks with various levels of accuracy. With some hardware the operating systems can pass the system clock time almost directly to programs. A system clock can also be derived from a chip that simply provides timing interrupts, periodic ticks at some known time interval. In either case, operating system kernels provide a clock source that hides these details. But the accuracy of that clock source and how quickly it can return results varies based on the underlying hardware.

Inaccurate time keeping can result in system instability. Test any change to the clock source very carefully. Operating system defaults are sometimes made to favor reliability over best accuracy. And if you are using a virtual machine, look into the recommended time sources compatible with it. Virtual hardware faces additional difficulties when emulating timers, and there are often per operating system settings suggested by vendors.

The Time Stamp Counter (TSC) clock source is the most accurate one available on current generation CPUs. It's the preferred way to track the system time when it's supported by the operating system and the TSC clock is reliable. There are several ways that TSC can fail to provide an accurate timing source, making it unreliable. Older systems can have a TSC clock that varies based on the CPU temperature, making it unusable for timing. Trying to use TSC on some older multicore CPUs can give a reported time that's inconsistent among multiple cores. This can result in the time going backwards, a problem this program checks for. And even the newest systems can fail to provide accurate TSC timing with very aggressive power saving configurations.

Newer operating systems may check for the known TSC problems and switch to a slower, more stable clock source when they are seen. If your system supports TSC time but doesn't default to that, it may be disabled for a good reason. And some operating systems may not detect all the possible problems correctly, or will allow using TSC even in situations where it's known to be inaccurate.

The High Precision Event Timer (HPET) is the preferred timer on systems where it's available and TSC is not accurate. The timer chip itself is programmable to allow up to 100 nanosecond resolution, but you may not see that much accuracy in your system clock.

Advanced Configuration and Power Interface (ACPI) provides a Power Management (PM) Timer, which Linux refers to as the acpi_pm. The clock derived from acpi_pm will at best provide 300 nanosecond resolution.
Timers used on older PC hardware include the 8254 Programmable Interval Timer (PIT), the real-time clock (RTC), the Advanced Programmable Interrupt Controller (APIC) timer, and the Cyclone timer. These timers aim for millisecond resolution.

**See Also**

EXPLAIN
pg_upgrade

pg_upgrade — upgrade a PostgreSQL server instance

Synopsis

pg_upgrade -b oldbindir -B newbindir -d oldconfigdir -D newconfigdir[option...]

Description

pg_upgrade (formerly called pg_migrator) allows data stored in PostgreSQL data files to be upgraded
to a later PostgreSQL major version without the data dump/reload typically required for major version
upgrades, e.g. from 9.5.8 to 9.6.4 or from 10.7 to 11.2. It is not required for minor version upgrades,
e.g. from 9.6.2 to 9.6.3 or from 10.1 to 10.2.

Major PostgreSQL releases regularly add new features that often change the layout of the system ta-
bles, but the internal data storage format rarely changes. pg_upgrade uses this fact to perform rapid
upgrades by creating new system tables and simply reusing the old user data files. If a future major re-
lease ever changes the data storage format in a way that makes the old data format unreadable, pg_up-
grade will not be usable for such upgrades. (The community will attempt to avoid such situations.)

pg_upgrade does its best to make sure the old and new clusters are binary-compatible, e.g. by checking
for compatible compile-time settings, including 32/64-bit binaries. It is important that any external
modules are also binary compatible, though this cannot be checked by pg_upgrade.

pg_upgrade supports upgrades from 8.4.X and later to the current major release of PostgreSQL, in-
cluding snapshot and beta releases.

Options

pg_upgrade accepts the following command-line arguments:

- \[b\] bindir
  --old-bindir=bindir
  the old PostgreSQL executable directory; environment variable PGBINOLD

- \[B\] bindir
  --new-bindir=bindir
  the new PostgreSQL executable directory; environment variable PGBINNEW

- \[c\]
  --check
  check clusters only, don’t change any data

- \[d\] configdir
  --old-datadir=configdir
  the old database cluster configuration directory; environment variable PGDATAOLD

- \[D\] configdir
  --new-datadir=configdir
  the new database cluster configuration directory; environment variable PGDATANEW
-j
--jobs

   number of simultaneous processes or threads to use

-k
--link

   use hard links instead of copying files to the new cluster

-o options
--old-options options

   options to be passed directly to the old postgres command; multiple option invocations are appended

-O options
--new-options options

   options to be passed directly to the new postgres command; multiple option invocations are appended

-p port
--old-port=port

   the old cluster port number; environment variable PGPORTOLD

-P port
--new-port=port

   the new cluster port number; environment variable PGPORTNEW

-r
--retain

   retain SQL and log files even after successful completion

-U username
--username=username

   cluster's install user name; environment variable PGUSER

-v
--verbose

   enable verbose internal logging

-V
--version

   display version information, then exit

-?
--help

   show help, then exit

**Usage**

These are the steps to perform an upgrade with pg_upgrade:
1. **Optionally move the old cluster**

If you are using a version-specific installation directory, e.g. /opt/PostgreSQL/11, you do not need to move the old cluster. The graphical installers all use version-specific installation directories.

If your installation directory is not version-specific, e.g. /usr/local/pgsql, it is necessary to move the current PostgreSQL install directory so it does not interfere with the new PostgreSQL installation. Once the current PostgreSQL server is shut down, it is safe to rename the PostgreSQL installation directory; assuming the old directory is /usr/local/pgsql, you can do:

```sh
cd /usr/local/pgsql
mv /usr/local/pgsql /usr/local/pgsql.old
```
to rename the directory.

2. **For source installs, build the new version**

Build the new PostgreSQL source with `configure` flags that are compatible with the old cluster. `pg_upgrade` will check `pg_controldata` to make sure all settings are compatible before starting the upgrade.

3. **Install the new PostgreSQL binaries**

Install the new server's binaries and support files. `pg_upgrade` is included in a default installation.

For source installs, if you wish to install the new server in a custom location, use the `prefix` variable:

```sh
make prefix=/usr/local/pgsql.new install
```

4. **Initialize the new PostgreSQL cluster**

Initialize the new cluster using `initdb`. Again, use compatible `initdb` flags that match the old cluster. Many prebuilt installers do this step automatically. There is no need to start the new cluster.

5. **Install custom shared object files**

Install any custom shared object files (or DLLs) used by the old cluster into the new cluster, e.g. `pgcrypto.so`, whether they are from `contrib` or some other source. Do not install the schema definitions, e.g. `CREATE EXTENSION pgcrypto`, because these will be upgraded from the old cluster. Also, any custom full text search files (dictionary, synonym, thesaurus, stop words) must also be copied to the new cluster.

6. **Adjust authentication**

`pg_upgrade` will connect to the old and new servers several times, so you might want to set authentication to `peer` in `pg_hba.conf` or use a `~/.pgpass` file (see Section 34.15).

7. **Stop both servers**

Make sure both database servers are stopped using, on Unix, e.g.:

```sh
pg_ctl -D /opt/PostgreSQL/9.6 stop
pg_ctl -D /opt/PostgreSQL/11 stop
```
or on Windows, using the proper service names:
NET STOP postgresql-9.6
NET STOP postgresql-11

Streaming replication and log-shipping standby servers can remain running until a later step.

8. **Prepare for standby server upgrades**

If you are upgrading standby servers using methods outlined in section Step 10, verify that the old standby servers are caught up by running `pg_controldata` against the old primary and standby clusters. Verify that the “Latest checkpoint location” values match in all clusters. (There will be a mismatch if old standby servers were shut down before the old primary or if the old standby servers are still running.) Also, change `wal_level` to `replica` in the `postgresql.conf` file on the new primary cluster.

9. **Run pg_upgrade**

Always run the `pg_upgrade` binary of the new server, not the old one. `pg_upgrade` requires the specification of the old and new cluster's data and executable (bin) directories. You can also specify user and port values, and whether you want the data files linked instead of the default copy behavior.

If you use link mode, the upgrade will be much faster (no file copying) and use less disk space, but you will not be able to access your old cluster once you start the new cluster after the upgrade. Link mode also requires that the old and new cluster data directories be in the same file system. (Tablespaces and `pg_wal` can be on different file systems.) See `pg_upgrade --help` for a full list of options.

The `--jobs` option allows multiple CPU cores to be used for copying/linking of files and to dump and reload database schemas in parallel; a good place to start is the maximum of the number of CPU cores and tablespaces. This option can dramatically reduce the time to upgrade a multi-database server running on a multiprocessor machine.

For Windows users, you must be logged into an administrative account, and then start a shell as the `postgres` user and set the proper path:

```
RUNAS /USER:postgres "CMD.EXE"
SET PATH=%PATH%;C:\Program Files\PostgreSQL\11\bin;
```

and then run `pg_upgrade` with quoted directories, e.g.:

```
pg_upgrade.exe
    --old-datadir "C:/Program Files/PostgreSQL/9.6/data"
    --new-datadir "C:/Program Files/PostgreSQL/11/data"
    --old-bindir "C:/Program Files/PostgreSQL/9.6/bin"
    --new-bindir "C:/Program Files/PostgreSQL/11/bin"
```

Once started, `pg_upgrade` will verify the two clusters are compatible and then do the upgrade. You can use `pg_upgrade --check` to perform only the checks, even if the old server is still running. `pg_upgrade --check` will also outline any manual adjustments you will need to make after the upgrade. If you are going to be using link mode, you should use the `--link` option with `--check` to enable link-mode-specific checks. `pg_upgrade` requires write permission in the current directory.

Obviously, no one should be accessing the clusters during the upgrade. `pg_upgrade` defaults to running servers on port 50432 to avoid unintended client connections. You can use the same port number for both clusters when doing an upgrade because the old and new clusters will not be running at the same time. However, when checking an old running server, the old and new port numbers must be different.
If an error occurs while restoring the database schema, `pg_upgrade` will exit and you will have to revert to the old cluster as outlined in Step 16 below. To try `pg_upgrade` again, you will need to modify the old cluster so the `pg_upgrade` schema restore succeeds. If the problem is a `contrib` module, you might need to uninstall the `contrib` module from the old cluster and install it in the new cluster after the upgrade, assuming the module is not being used to store user data.

10. **Upgrade Streaming Replication and Log-Shipping standby servers**

If you used link mode and have Streaming Replication (see Section 26.2.5) or Log-Shipping (see Section 26.2) standby servers, you can follow these steps to quickly upgrade them. You will not be running `pg_upgrade` on the standby servers, but rather `rsync` on the primary. Do not start any servers yet.

If you did *not* use link mode, do not have or do not want to use `rsync`, or want an easier solution, skip the instructions in this section and simply recreate the standby servers once `pg_upgrade` completes and the new primary is running.

a. **Install the new PostgreSQL binaries on standby servers**

   Make sure the new binaries and support files are installed on all standby servers.

b. **Make sure the new standby data directories do not exist**

   Make sure the new standby data directories do not exist or are empty. If `initdb` was run, delete the standby servers' new data directories.

c. **Install custom shared object files**

   Install the same custom shared object files on the new standbys that you installed in the new primary cluster.

d. **Stop standby servers**

   If the standby servers are still running, stop them now using the above instructions.

e. **Save configuration files**

   Save any configuration files from the old standbys' configuration directories you need to keep, e.g. `postgresql.conf`, `recovery.conf`, because these will be overwritten or removed in the next step.

f. **Run rsync**

   When using link mode, standby servers can be quickly upgraded using `rsync`. To accomplish this, from a directory on the primary server that is above the old and new database cluster directories, run this on the `primary` for each standby server:

   ```bash
   rsync --archive --delete --hard-links --size-only --no-inc-recursive old_cluster new_cluster remote_dir
   ```

   where `old_cluster` and `new_cluster` are relative to the current directory on the primary, and `remote_dir` is above the old and new cluster directories on the standby. The directory structure under the specified directories on the primary and standbys must match. Consult the `rsync` manual page for details on specifying the remote directory, e.g.

   ```bash
   rsync --archive --delete --hard-links --size-only --no-inc-recursive /opt/PostgreSQL/9.5 /
   ```
You can verify what the command will do using rsync's --dry-run option. While rsync must be run on the primary for at least one standby, it is possible to run rsync on an upgraded standby to upgrade other standbys, as long as the upgraded standby has not been started.

What this does is to record the links created by pg_upgrade's link mode that connect files in the old and new clusters on the primary server. It then finds matching files in the standby's old cluster and creates links for them in the standby's new cluster. Files that were not linked on the primary are copied from the primary to the standby. (They are usually small.) This provides rapid standby upgrades. Unfortunately, rsync needlessly copies files associated with temporary and unlogged tables because these files don't normally exist on standby servers.

If you have tablespaces, you will need to run a similar rsync command for each tablespace directory, e.g.:

```
rsync --archive --delete --hard-links --size-only --no-inc-recursive /vol1/pg_tblsp/PG_9.5_201510051 /vol1/pg_tblsp/PG_9.6_201608131 standby.example.com:/vol1/pg_tblsp
```

If you have relocated pg_wal outside the data directories, rsync must be run on those directories too.

g. **Configure streaming replication and log-shipping standby servers**

Configure the servers for log shipping. (You do not need to run pg_start_backup() and pg_stop_backup() or take a file system backup as the standbys are still synchronized with the primary.)

11. **Restore pg_hba.conf**

If you modified pg_hba.conf, restore its original settings. It might also be necessary to adjust other configuration files in the new cluster to match the old cluster, e.g. postgresql.conf.

12. **Start the new server**

The new server can now be safely started, and then any rsync'ed standby servers.

13. **Post-Upgrade processing**

If any post-upgrade processing is required, pg_upgrade will issue warnings as it completes. It will also generate script files that must be run by the administrator. The script files will connect to each database that needs post-upgrade processing. Each script should be run using:

```
psql --username=postgres --file=script.sql postgres
```

The scripts can be run in any order and can be deleted once they have been run.

---

**Caution**

In general it is unsafe to access tables referenced in rebuild scripts until the rebuild scripts have run to completion; doing so could yield incorrect results or poor performance. Tables not referenced in rebuild scripts can be accessed immediately.
14. **Statistics**

Because optimizer statistics are not transferred by `pg_upgrade`, you will be instructed to run a command to regenerate that information at the end of the upgrade. You might need to set connection parameters to match your new cluster.

15. **Delete old cluster**

Once you are satisfied with the upgrade, you can delete the old cluster's data directories by running the script mentioned when `pg_upgrade` completes. (Automatic deletion is not possible if you have user-defined tablespaces inside the old data directory.) You can also delete the old installation directories (e.g., `bin`, `share`).

16. **Reverting to old cluster**

If, after running `pg_upgrade`, you wish to revert to the old cluster, there are several options:

- If you ran `pg_upgrade` with `--check`, no modifications were made to the old cluster and you can re-use it anytime.

- If you ran `pg_upgrade` with `--link`, the data files are shared between the old and new cluster. If you started the new cluster, the new server has written to those shared files and it is unsafe to use the old cluster.

- If you ran `pg_upgrade` without `--link` or did not start the new server, the old cluster was not modified except that, if linking started, a `.old` suffix was appended to `$PGDATA/global/pg_control`. To reuse the old cluster, possibly remove the `.old` suffix from `$PGDATA/global/pg_control`; you can then restart the old cluster.

**Notes**

`pg_upgrade` does not support upgrading of databases containing table columns using these `reg* OID-referencing system data types: regproc, regprocedure, regoper, regoperator, regconfig, and regdictionary. (regtype can be upgraded.)`

All failure, rebuild, and reindex cases will be reported by `pg_upgrade` if they affect your installation; post-upgrade scripts to rebuild tables and indexes will be generated automatically. If you are trying to automate the upgrade of many clusters, you should find that clusters with identical database schemas require the same post-upgrade steps for all cluster upgrades; this is because the post-upgrade steps are based on the database schemas, and not user data.

For deployment testing, create a schema-only copy of the old cluster, insert dummy data, and upgrade that.

If you are upgrading a pre-PostgreSQL 9.2 cluster that uses a configuration-file-only directory, you must pass the real data directory location to `pg_upgrade`, and pass the configuration directory location to the server, e.g., `-d /real-data-directory -o '-D /configuration-directory'`.

If using a pre-9.1 old server that is using a non-default Unix-domain socket directory or a default that differs from the default of the new cluster, set `PGHOST` to point to the old server's socket location. (This is not relevant on Windows.)

If you want to use link mode and you do not want your old cluster to be modified when the new cluster is started, make a copy of the old cluster and upgrade that in link mode. To make a valid copy of the old cluster, use `rsync` to create a dirty copy of the old cluster while the server is running, then shut down the old server and run `rsync --checksum` again to update the copy with any changes to make it consistent. (`--checksum` is necessary because `rsync` only has file modification-time granularity of one second.) You might want to exclude some files, e.g., `postmaster.pid`, as documented in Section 25.3.3. If your file system supports file system snapshots or copy-on-write file copies, you
can use that to make a backup of the old cluster and tablespaces, though the snapshot and copies must be created simultaneously or while the database server is down.

See Also
initdb, pg_ctl, pg_dump, postgres
pg_verify_checksums

pg_verify_checksums — verify data checksums in a PostgreSQL database cluster

Synopsis

pg_verify_checksums [option...] [-D --pgdata] datadir

Description

pg_verify_checksums verifies data checksums in a PostgreSQL cluster. The server must be shut down cleanly before running pg_verify_checksums. The exit status is zero if there are no checksum errors, otherwise nonzero.

Options

The following command-line options are available:

- -D directory
  --pgdata=directory

  Specifies the directory where the database cluster is stored.

- -v
  --verbose

  Enable verbose output. Lists all checked files.

- -r relfilenode

  Only validate checksums in the relation with specified relfilenode.

- -V
  --version

  Print the pg_verify_checksums version and exit.

- -?
  --help

  Show help about pg_verify_checksums command line arguments, and exit.

Environment

PGDATA

  Specifies the directory where the database cluster is stored; can be overridden using the -D option.
pg_waldump

pg_waldump — display a human-readable rendering of the write-ahead log of a PostgreSQL database cluster

Synopsis

pg_waldump [option...] [startseg [endseg]]

Description

pg_waldump displays the write-ahead log (WAL) and is mainly useful for debugging or educational purposes.

This utility can only be run by the user who installed the server, because it requires read-only access to the data directory.

Options

The following command-line options control the location and format of the output:

startseg

Start reading at the specified log segment file. This implicitly determines the path in which files will be searched for, and the timeline to use.

endseg

Stop after reading the specified log segment file.

-b
--b kp-details

Output detailed information about backup blocks.

-e end
--end=end

Stop reading at the specified WAL location, instead of reading to the end of the log stream.

-f
--follow

After reaching the end of valid WAL, keep polling once per second for new WAL to appear.

-n limit
--limit=limit

Display the specified number of records, then stop.

-p path
--path=path

Specifies a directory to search for log segment files or a directory with a pg_wal subdirectory that contains such files. The default is to search in the current directory, the pg_wal subdirectory of the current directory, and the pg_wal subdirectory of PGDATA.
pg_waldump

-x rmgr
--rmgr=rmgr

Only display records generated by the specified resource manager. If list is passed as name, print a list of valid resource manager names, and exit.

-s start
--start=start

WAL location at which to start reading. The default is to start reading the first valid log record found in the earliest file found.

-t timeline
--timeline=timeline

Timeline from which to read log records. The default is to use the value in startseg, if that is specified; otherwise, the default is 1.

-V
--version

Print the pg_waldump version and exit.

-x xid
--xid=xid

Only display records marked with the given transaction ID.

-z
--stats[=record]

Display summary statistics (number and size of records and full-page images) instead of individual records. Optionally generate statistics per-record instead of per-rmgr.

-?
--help

Show help about pg_waldump command line arguments, and exit.

Notes

Can give wrong results when the server is running.

Only the specified timeline is displayed (or the default, if none is specified). Records in other timelines are ignored.

pg_waldump cannot read WAL files with suffix .partial. If those files need to be read, .partial suffix needs to be removed from the file name.

See Also

Section 30.5
**postgres**

postgres — PostgreSQL database server

**Synopsis**

    postgres [option...]

**Description**

`postgres` is the PostgreSQL database server. In order for a client application to access a database it connects (over a network or locally) to a running `postgres` instance. The `postgres` instance then starts a separate server process to handle the connection.

One `postgres` instance always manages the data of exactly one database cluster. A database cluster is a collection of databases that is stored at a common file system location (the “data area”). More than one `postgres` instance can run on a system at one time, so long as they use different data areas and different communication ports (see below). When `postgres` starts it needs to know the location of the data area. The location must be specified by the `-D` option or the `PGDATA` environment variable; there is no default. Typically, `-D` or `PGDATA` points directly to the data area directory created by `initdb`. Other possible file layouts are discussed in Section 19.2.

By default `postgres` starts in the foreground and prints log messages to the standard error stream. In practical applications `postgres` should be started as a background process, perhaps at boot time.

The `postgres` command can also be called in single-user mode. The primary use for this mode is during bootstrapping by `initdb`. Sometimes it is used for debugging or disaster recovery: note that running a single-user server is not truly suitable for debugging the server, since no realistic interprocess communication and locking will happen. When invoked in single-user mode from the shell, the user can enter queries and the results will be printed to the screen, but in a form that is more useful for developers than end users. In the single-user mode, the session user will be set to the user with ID 1, and implicit superuser powers are granted to this user. This user does not actually have to exist, so the single-user mode can be used to manually recover from certain kinds of accidental damage to the system catalogs.

**Options**

`postgres` accepts the following command-line arguments. For a detailed discussion of the options consult Chapter 19. You can save typing most of these options by setting up a configuration file. Some (safe) options can also be set from the connecting client in an application-dependent way to apply only for that session. For example, if the environment variable `PGOPTIONS` is set, then libpq-based clients will pass that string to the server, which will interpret it as `postgres` command-line options.

**General Purpose**

- `-B nbuffers`
  
  Sets the number of shared buffers for use by the server processes. The default value of this parameter is chosen automatically by `initdb`. Specifying this option is equivalent to setting the `shared_buffers` configuration parameter.

- `-c name=value`
  
  Sets a named run-time parameter. The configuration parameters supported by PostgreSQL are described in Chapter 19. Most of the other command line options are in fact short forms of such a parameter assignment. `-c` can appear multiple times to set multiple parameters.
-C name

Prints the value of the named run-time parameter, and exits. (See the -c option above for details.) This can be used on a running server, and returns values from postgresql.conf, modified by any parameters supplied in this invocation. It does not reflect parameters supplied when the cluster was started.

This option is meant for other programs that interact with a server instance, such as pg_ctl, to query configuration parameter values. User-facing applications should instead use SHOW or the pg_settings view.

-d debug-level

Sets the debug level. The higher this value is set, the more debugging output is written to the server log. Values are from 1 to 5. It is also possible to pass -d 0 for a specific session, which will prevent the server log level of the parent postgres process from being propagated to this session.

-D datadir

Specifies the file system location of the database configuration files. See Section 19.2 for details.

-e

Sets the default date style to “European”, that is DMY ordering of input date fields. This also causes the day to be printed before the month in certain date output formats. See Section 8.5 for more information.

-F

Disables fsync calls for improved performance, at the risk of data corruption in the event of a system crash. Specifying this option is equivalent to disabling the fsync configuration parameter. Read the detailed documentation before using this!

-h hostname

Specifies the IP host name or address on which postgres is to listen for TCP/IP connections from client applications. The value can also be a comma-separated list of addresses, or * to specify listening on all available interfaces. An empty value specifies not listening on any IP addresses, in which case only Unix-domain sockets can be used to connect to the server. Defaults to listening only on localhost. Specifying this option is equivalent to setting the listen_addresses configuration parameter.

-i

Allows remote clients to connect via TCP/IP (Internet domain) connections. Without this option, only local connections are accepted. This option is equivalent to setting listen_addresses to * in postgresql.conf or via -h.

This option is deprecated since it does not allow access to the full functionality of listen_addresses. It's usually better to set listen_addresses directly.

-k directory

Specifies the directory of the Unix-domain socket on which postgres is to listen for connections from client applications. The value can also be a comma-separated list of directories. An empty value specifies not listening on any Unix-domain sockets, in which case only TCP/IP sockets can be used to connect to the server. The default value is normally /tmp, but that can be changed at build time. Specifying this option is equivalent to setting the unix_socket_directories configuration parameter.
-l
Enables secure connections using SSL. PostgreSQL must have been compiled with support for
SSL for this option to be available. For more information on using SSL, refer to Section 18.9.

-N max-connections
Sets the maximum number of client connections that this server will accept. The default value of
this parameter is chosen automatically by initdb. Specifying this option is equivalent to setting
the max_connections configuration parameter.

-o extra-options
The command-line-style arguments specified in extra-options are passed to all server
processes started by this postgres process.

Spaces within extra-options are considered to separate arguments, unless escaped with a
backslash (\); write \ \ to represent a literal backslash. Multiple arguments can also be specified
via multiple uses of -o.

The use of this option is obsolete; all command-line options for server processes can be specified
directly on the postgres command line.

-p port
Specifies the TCP/IP port or local Unix domain socket file extension on which
postgres is to
listen for connections from client applications. Defaults to the value of the PGPORT environment
variable, or if PGPORT is not set, then defaults to the value established during compilation (nor-
manly 5432). If you specify a port other than the default port, then all client applications must
specify the same port using either command-line options or PGPORT.

-s
Print time information and other statistics at the end of each command. This is useful for bench-
marking or for use in tuning the number of buffers.

-S work-mem
Specifies the amount of memory to be used by internal sorts and hashes before resorting to tem-
porary disk files. See the description of the work_mem configuration parameter in Section 19.4.1.

-V
--version
Print the postgres version and exit.

--name=value
Sets a named run-time parameter; a shorter form of -c.

--describe-config
This option dumps out the server's internal configuration variables, descriptions, and defaults in
tab-delimited COPY format. It is designed primarily for use by administration tools.

-?
--help
Show help about postgres command line arguments, and exit.

Semi-internal Options
The options described here are used mainly for debugging purposes, and in some cases to assist with
recovery of severely damaged databases. There should be no reason to use them in a production data-
base setup. They are listed here only for use by PostgreSQL system developers. Furthermore, these options might change or be removed in a future release without notice.

\[-f \{ s | i | o | b | t | n | m | h \}\]

Forbids the use of particular scan and join methods: \(s\) and \(i\) disable sequential and index scans respectively, \(o\), \(b\) and \(t\) disable index-only scans, bitmap index scans, and TID scans respectively, while \(n\), \(m\), and \(h\) disable nested-loop, merge and hash joins respectively.

Neither sequential scans nor nested-loop joins can be disabled completely; the \(-fs\) and \(-fn\) options simply discourage the optimizer from using those plan types if it has any other alternative.

\(-n\)

This option is for debugging problems that cause a server process to die abnormally. The ordinary strategy in this situation is to notify all other server processes that they must terminate and then reinitialize the shared memory and semaphores. This is because an errant server process could have corrupted some shared state before terminating. This option specifies that \(postgres\) will not reinitialize shared data structures. A knowledgeable system programmer can then use a debugger to examine shared memory and semaphore state.

\(-O\)

Allows the structure of system tables to be modified. This is used by \(initdb\).

\(-P\)

Ignore system indexes when reading system tables, but still update the indexes when modifying the tables. This is useful when recovering from damaged system indexes.

\(-t pa[rser] | pl[anner] | e[xecutor]\)

Print timing statistics for each query relating to each of the major system modules. This option cannot be used together with the \(-s\) option.

\(-T\)

This option is for debugging problems that cause a server process to die abnormally. The ordinary strategy in this situation is to notify all other server processes that they must terminate and then reinitialize the shared memory and semaphores. This is because an errant server process could have corrupted some shared state before terminating. This option specifies that \(postgres\) will stop all other server processes by sending the signal \(SIGSTOP\), but will not cause them to terminate. This permits system programmers to collect core dumps from all server processes by hand.

\(-v protocol\)

Specifies the version number of the frontend/backend protocol to be used for a particular session. This option is for internal use only.

\(-W seconds\)

A delay of this many seconds occurs when a new server process is started, after it conducts the authentication procedure. This is intended to give an opportunity to attach to the server process with a debugger.

**Options for Single-User Mode**

The following options only apply to the single-user mode (see Single-User Mode).

\(--single\)

Selects the single-user mode. This must be the first argument on the command line.
**database**

Specifies the name of the database to be accessed. This must be the last argument on the command line. If it is omitted it defaults to the user name.

- **E**

Echo all commands to standard output before executing them.

- **j**

Use semicolon followed by two newlines, rather than just newline, as the command entry terminator.

- **r filename**

Send all server log output to `filename`. This option is only honored when supplied as a command-line option.

**Environment**

- **PGCLIENTENCODING**

Default character encoding used by clients. (The clients can override this individually.) This value can also be set in the configuration file.

- **PGDATA**

Default data directory location

- **PGDATESTYLE**

Default value of the DateStyle run-time parameter. (The use of this environment variable is deprecated.)

- **PGPORT**

Default port number (preferably set in the configuration file)

**Diagnostics**

A failure message mentioning `semget` or `shmget` probably indicates you need to configure your kernel to provide adequate shared memory and semaphores. For more discussion see Section 18.4. You might be able to postpone reconfiguring your kernel by decreasing `shared_buffers` to reduce the shared memory consumption of PostgreSQL, and/or by reducing `max_connections` to reduce the semaphore consumption.

A failure message suggesting that another server is already running should be checked carefully, for example by using the command

```
$ ps ax | grep postgres
```

or

```
$ ps -ef | grep postgres
```

depending on your system. If you are certain that no conflicting server is running, you can remove the lock file mentioned in the message and try again.
A failure message indicating inability to bind to a port might indicate that that port is already in use by some non-PostgreSQL process. You might also get this error if you terminate `postgres` and immediately restart it using the same port; in this case, you must simply wait a few seconds until the operating system closes the port before trying again. Finally, you might get this error if you specify a port number that your operating system considers to be reserved. For example, many versions of Unix consider port numbers under 1024 to be “trusted” and only permit the Unix superuser to access them.

**Notes**

The utility command `pg_ctl` can be used to start and shut down the `postgres` server safely and comfortably.

If at all possible, do not use `SIGKILL` to kill the main `postgres` server. Doing so will prevent `postgres` from freeing the system resources (e.g., shared memory and semaphores) that it holds before terminating. This might cause problems for starting a fresh `postgres` run.

To terminate the `postgres` server normally, the signals `SIGTERM`, `SIGINT`, or `SIGQUIT` can be used. The first will wait for all clients to terminate before quitting, the second will forcefully disconnect all clients, and the third will quit immediately without proper shutdown, resulting in a recovery run during restart.

The `SIGHUP` signal will reload the server configuration files. It is also possible to send `SIGHUP` to an individual server process, but that is usually not sensible.

To cancel a running query, send the `SIGINT` signal to the process running that command. To terminate a backend process cleanly, send `SIGTERM` to that process. See also `pg_cancel_backend` and `pg_terminate_backend` in Section 9.26.2 for the SQL-callable equivalents of these two actions.

The `postgres` server uses `SIGQUIT` to tell subordinate server processes to terminate without normal cleanup. This signal should not be used by users. It is also unwise to send `SIGKILL` to a server process — the main `postgres` process will interpret this as a crash and will force all the sibling processes to quit as part of its standard crash-recovery procedure.

**Bugs**

The `--` options will not work on FreeBSD or OpenBSD. Use `–c` instead. This is a bug in the affected operating systems; a future release of PostgreSQL will provide a workaround if this is not fixed.

**Single-User Mode**

To start a single-user mode server, use a command like

```bash
postgres --single -D /usr/local/pgsql/data other-options my_database
```

Provide the correct path to the database directory with `–D`, or make sure that the environment variable `PGDATA` is set. Also specify the name of the particular database you want to work in.

Normally, the single-user mode server treats newline as the command entry terminator; there is no intelligence about semicolons, as there is in `psql`. To continue a command across multiple lines, you must type backslash just before each newline except the last one. The backslash and adjacent newline are both dropped from the input command. Note that this will happen even when within a string literal or comment.

But if you use the `–j` command line switch, a single newline does not terminate command entry; instead, the sequence semicolon-newline-newline does. That is, type a semicolon immediately followed by a completely empty line. Backslash-newline is not treated specially in this mode. Again, there is no intelligence about such a sequence appearing within a string literal or comment.
In either input mode, if you type a semicolon that is not just before or part of a command entry terminator, it is considered a command separator. When you do type a command entry terminator, the multiple statements you've entered will be executed as a single transaction.

To quit the session, type EOF (Control+D, usually). If you've entered any text since the last command entry terminator, then EOF will be taken as a command entry terminator, and another EOF will be needed to exit.

Note that the single-user mode server does not provide sophisticated line-editing features (no command history, for example). Single-user mode also does not do any background processing, such as automatic checkpoints or replication.

Examples

To start `postgres` in the background using default values, type:

```
$ nohup postgres >logfile 2>&1 </dev/null &
```

To start `postgres` with a specific port, e.g. 1234:

```
$ postgres -p 1234
```

To connect to this server using `psql`, specify this port with the `-p` option:

```
$ psql -p 1234
```

or set the environment variable `PGPORT`:

```
$ export PGPORT=1234
$ psql
```

Named run-time parameters can be set in either of these styles:

```
$ postgres -c work_mem=1234
$ postgres --work-mem=1234
```

Either form overrides whatever setting might exist for `work_mem` in `postgresql.conf`. Notice that underscores in parameter names can be written as either underscore or dash on the command line. Except for short-term experiments, it's probably better practice to edit the setting in `postgresql.conf` than to rely on a command-line switch to set a parameter.

See Also

`initdb`, `pg_ctl`
postmaster

postmaster — PostgreSQL database server

Synopsis

postmaster [option...]

Description

postmaster is a deprecated alias of postgres.

See Also

postgres
Part VII. Internals

This part contains assorted information that might be of use to PostgreSQL developers.
## Table of Contents

51. Overview of PostgreSQL Internals ................................................................. 1994
   51.1. The Path of a Query ........................................................................ 1994
   51.2. How Connections are Established .................................................. 1994
   51.3. The Parser Stage ............................................................................ 1995
      51.3.1. Parser .................................................................................... 1995
      51.3.2. Transformation Process .......................................................... 1995
   51.4. The PostgreSQL Rule System ............................................................ 1996
   51.5. Planner/Optimizer ........................................................................... 1996
      51.5.1. Generating Possible Plans ....................................................... 1997
   51.6. Executor ......................................................................................... 1998

52. System Catalogs ....................................................................................... 1999
   52.1. Overview ........................................................................................... 1999
   52.2. pg_aggregate .................................................................................... 2000
   52.3. pg_am .............................................................................................. 2003
   52.4. pg_amop .......................................................................................... 2003
   52.5. pg_amproc ...................................................................................... 2004
   52.6. pg_attrdef ...................................................................................... 2005
   52.7. pg_attribute ..................................................................................... 2005
   52.8. pg_authid ........................................................................................ 2009
   52.9. pg_auth_members ............................................................................. 2010
   52.10. pg_cast ............................................................................................ 2010
   52.11. pg_class .......................................................................................... 2011
   52.12. pg_collation ................................................................................... 2015
   52.13. pg_constraint .................................................................................. 2016
   52.14. pg_conversion ................................................................................ 2019
   52.15. pg_database .................................................................................... 2019
   52.16. pg_db_role_setting ....................................................................... 2021
   52.17. pg_default_acl ............................................................................. 2022
   52.18. pg_depend ...................................................................................... 2022
   52.19. pg_description ............................................................................. 2024
   52.20. pg_enum .......................................................................................... 2024
   52.21. pg_event_trigger .......................................................................... 2025
   52.22. pg_extension .................................................................................. 2025
   52.23. pg_foreign_data_wrapper .............................................................. 2026
   52.24. pg_foreign_server ......................................................................... 2026
   52.25. pg_foreign_table .......................................................................... 2027
   52.26. pg_index ......................................................................................... 2028
   52.27. pg_inherits ..................................................................................... 2028
   52.28. pg_init_privs .................................................................................. 2031
   52.29. pg_language ................................................................................... 2031
   52.30. pg_largeobject .............................................................................. 2032
   52.31. pg_largeobject_metadata ............................................................... 2033
   52.32. pg_namespace .................................................................................. 2034
   52.33. pg_opclass ..................................................................................... 2034
   52.34. pg_operator .................................................................................... 2035
   52.35. pg_opfamily ................................................................................... 2036
   52.36. pg_partitioned_table ................................................................. 2036
   52.37. pg_pltemplate .............................................................................. 2038
   52.38. pg_policy ....................................................................................... 2038
   52.39. pg_proc .......................................................................................... 2039
   52.40. pg_publication ............................................................................. 2043
   52.41. pg_publication_rel ....................................................................... 2044
   52.42. pg_range ....................................................................................... 2044
   52.43. pg_replication_origin ................................................................. 2045
   52.44. pg_rewrite ..................................................................................... 2045
52.45. pg_seclabel ................................................................. 2046
52.46. pg_sequence ............................................................. 2046
52.47. pg_shdepend ............................................................ 2047
52.48. pg_shdescription ...................................................... 2048
52.49. pg_shseclabel .......................................................... 2049
52.50. pg_statistic ............................................................ 2049
52.51. pg_statistic_ext ........................................................ 2051
52.52. pg_subscription ....................................................... 2052
52.53. pg_subscription_rel .................................................. 2053
52.54. pg_tablespace .......................................................... 2053
52.55. pg_transform .......................................................... 2054
52.56. pg_trigger .............................................................. 2054
52.57. pg_ts_config ............................................................ 2056
52.58. pg_ts_config_map ...................................................... 2056
52.59. pg_ts_dict .............................................................. 2057
52.60. pg_ts_parser ............................................................ 2057
52.61. pg_ts_template ........................................................ 2058
52.62. pg_type ................................................................. 2058
52.63. pg_user_mapping ....................................................... 2065
52.64. System Views .......................................................... 2065
52.65. pg_available_extensions ............................................ 2066
52.66. pg_available_extension_versions .................................. 2067
52.67. pg_config .............................................................. 2067
52.68. pg_cursors ............................................................. 2068
52.69. pg_file_settings ...................................................... 2068
52.70. pg_group ............................................................... 2069
52.71. pg_hba_file_rules ..................................................... 2069
52.72. pg_indexes ............................................................. 2070
52.73. pg_locks ............................................................... 2071
52.74. pg_matviews ........................................................... 2073
52.75. pg_policies ............................................................ 2074
52.76. pg_prepared_statements ............................................. 2075
52.77. pg_prepared_xacts .................................................... 2075
52.78. pg_publication_tables .............................................. 2076
52.79. pg_replication_origin_status ..................................... 2076
52.80. pg_replication_slots ............................................... 2077
52.81. pg_roles ............................................................... 2078
52.82. pg_rules .............................................................. 2079
52.83. pg_seclabels .......................................................... 2080
52.84. pg_sequences ........................................................ 2080
52.85. pg_settings .......................................................... 2081
52.86. pg_shadow ............................................................ 2083
52.87. pg_stats ............................................................... 2084
52.88. pg_tables ............................................................. 2086
52.89. pg_timezone_abbrs ................................................... 2087
52.90. pg_timezone_names .................................................. 2087
52.91. pg_user ............................................................... 2088
52.92. pg_user_mappings .................................................... 2088
52.93. pg_views ............................................................. 2089
53. Frontend/Backend Protocol .............................................. 2090
53.1. Overview ............................................................... 2090
53.1.1. Messaging Overview .............................................. 2090
53.1.2. Extended Query Overview ...................................... 2091
53.1.3. Formats and Format Codes ..................................... 2091
53.2. Message Flow ........................................................ 2091
53.2.1. Start-up ............................................................ 2092
53.2.2. Simple Query .................................................... 2094
53.2.3. Extended Query .................................................. 2097
53.2.4. Function Call ................................................................. 2100
53.2.5. COPY Operations ....................................................... 2100
53.2.6. Asynchronous Operations .......................................... 2101
53.2.7. Canceling Requests in Progress .................................. 2102
53.2.8. Termination ............................................................... 2103
53.2.9. SSL Session Encryption ............................................. 2103
53.3. SASL Authentication ...................................................... 2103
53.4. Streaming Replication Protocol ....................................... 2105
53.5. Logical Streaming Replication Protocol ......................... 2112
  53.5.1. Logical Streaming Replication Parameters .................. 2112
  53.5.2. Logical Replication Protocol Messages ....................... 2112
  53.5.3. Logical Replication Protocol Message Flow ................. 2113
53.6. Message Data Types ..................................................... 2113
53.7. Message Formats ........................................................ 2113
53.8. Error and Notice Message Fields ................................... 2130
53.9. Logical Replication Message Formats ........................... 2132
53.10. Summary of Changes since Protocol 2.0 ....................... 2136
54. PostgreSQL Coding Conventions ..................................... 2138
  54.1. Formatting ............................................................. 2138
  54.2. Reporting Errors Within the Server .............................. 2138
  54.3. Error Message Style Guide ........................................ 2141
  54.4. Miscellaneous Coding Conventions ............................... 2145
55. Native Language Support ................................................. 2148
  55.1. For the Translator ..................................................... 2148
    55.1.1. Requirements .................................................... 2148
    55.1.2. Concepts ........................................................ 2148
    55.1.3. Creating and Maintaining Message Catalogs ............. 2149
    55.1.4. Editing the PO Files ......................................... 2150
  55.2. For the Programmer .................................................. 2150
    55.2.1. Mechanics ....................................................... 2150
    55.2.2. Message-writing Guidelines ................................. 2152
56. Writing A Procedural Language Handler .......................... 2154
57. Writing A Foreign Data Wrapper ...................................... 2157
  57.1. Foreign Data Wrapper Functions .................................. 2157
  57.2. Foreign Data Wrapper Callback Routines ....................... 2157
    57.2.1. FDW Routines For Scanning Foreign Tables ............. 2157
    57.2.2. FDW Routines For Scanning Foreign Joins ............... 2159
    57.2.3. FDW Routines For Planning Post-Scan/Join Processing .. 2160
    57.2.4. FDW Routines For Updating Foreign Tables ............. 2161
    57.2.5. FDW Routines For Row Locking ............................... 2166
    57.2.6. FDW Routines for EXPLAIN ................................... 2167
    57.2.7. FDW Routines for ANALYZE .................................. 2168
    57.2.8. FDW Routines For IMPORT FOREIGN SCHEMA .............. 2168
    57.2.9. FDW Routines for Parallel Execution ..................... 2169
    57.2.10. FDW Routines For reparameterization of paths ....... 2170
  57.3. Foreign Data Wrapper Helper Functions ......................... 2170
  57.4. Foreign Data Wrapper Query Planning ........................... 2171
  57.5. Row Locking in Foreign Data Wrappers ......................... 2174
58. Writing A Table Sampling Method ................................... 2176
  58.1. Sampling Method Support Functions ............................. 2176
59. Writing A Custom Scan Provider ..................................... 2179
  59.1. Creating Custom Scan Paths ....................................... 2179
    59.1.1. Custom Scan Path Callbacks ................................ 2180
  59.2. Creating Custom Scan Plans ....................................... 2180
    59.2.1. Custom Scan Plan Callbacks ................................ 2181
  59.3. Executing Custom Scans ............................................ 2181
    59.3.1. Custom Scan Execution Callbacks .......................... 2181
## Internals

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>60. Genetic Query Optimizer</td>
<td>2184</td>
</tr>
<tr>
<td>60.1. Query Handling as a Complex Optimization Problem</td>
<td>2184</td>
</tr>
<tr>
<td>60.2. Genetic Algorithms</td>
<td>2184</td>
</tr>
<tr>
<td>60.3. Genetic Query Optimization (GEQO) in PostgreSQL</td>
<td>2185</td>
</tr>
<tr>
<td>60.3.1. Generating Possible Plans with GEQO</td>
<td>2186</td>
</tr>
<tr>
<td>60.3.2. Future Implementation Tasks for PostgreSQL GEQO</td>
<td>2186</td>
</tr>
<tr>
<td>60.4. Further Reading</td>
<td>2186</td>
</tr>
<tr>
<td>61. Index Access Method Interface Definition</td>
<td>2188</td>
</tr>
<tr>
<td>61.1. Basic API Structure for Indexes</td>
<td>2188</td>
</tr>
<tr>
<td>61.2. Index Access Method Functions</td>
<td>2190</td>
</tr>
<tr>
<td>61.3. Index Scanning</td>
<td>2195</td>
</tr>
<tr>
<td>61.4. Index Locking Considerations</td>
<td>2197</td>
</tr>
<tr>
<td>61.5. Index Uniqueness Checks</td>
<td>2198</td>
</tr>
<tr>
<td>61.6. Index Cost Estimation Functions</td>
<td>2199</td>
</tr>
<tr>
<td>62. Generic WAL Records</td>
<td>2202</td>
</tr>
<tr>
<td>63. B-Tree Indexes</td>
<td>2204</td>
</tr>
<tr>
<td>63.1. Introduction</td>
<td>2204</td>
</tr>
<tr>
<td>63.2. Behavior of B-Tree Operator Classes</td>
<td>2204</td>
</tr>
<tr>
<td>63.3. B-Tree Support Functions</td>
<td>2205</td>
</tr>
<tr>
<td>63.4. Implementation</td>
<td>2206</td>
</tr>
<tr>
<td>64. GiST Indexes</td>
<td>2207</td>
</tr>
<tr>
<td>64.1. Introduction</td>
<td>2207</td>
</tr>
<tr>
<td>64.2. Built-in Operator Classes</td>
<td>2207</td>
</tr>
<tr>
<td>64.3. Extensibility</td>
<td>2208</td>
</tr>
<tr>
<td>64.4. Implementation</td>
<td>2217</td>
</tr>
<tr>
<td>64.4.1. GiST buffering build</td>
<td>2217</td>
</tr>
<tr>
<td>64.5. Examples</td>
<td>2217</td>
</tr>
<tr>
<td>65. SP-GiST Indexes</td>
<td>2219</td>
</tr>
<tr>
<td>65.1. Introduction</td>
<td>2219</td>
</tr>
<tr>
<td>65.2. Built-in Operator Classes</td>
<td>2219</td>
</tr>
<tr>
<td>65.3. Extensibility</td>
<td>2219</td>
</tr>
<tr>
<td>65.4. Implementation</td>
<td>2228</td>
</tr>
<tr>
<td>65.4.1. SP-GiST Limits</td>
<td>2228</td>
</tr>
<tr>
<td>65.4.2. SP-GiST Without Node Labels</td>
<td>2228</td>
</tr>
<tr>
<td>65.4.3. “All-the-same” Inner Tuples</td>
<td>2228</td>
</tr>
<tr>
<td>65.5. Examples</td>
<td>2229</td>
</tr>
<tr>
<td>66. GIN Indexes</td>
<td>2230</td>
</tr>
<tr>
<td>66.1. Introduction</td>
<td>2230</td>
</tr>
<tr>
<td>66.2. Built-in Operator Classes</td>
<td>2230</td>
</tr>
<tr>
<td>66.3. Extensibility</td>
<td>2230</td>
</tr>
<tr>
<td>66.4. Implementation</td>
<td>2233</td>
</tr>
<tr>
<td>66.4.1. GIN Fast Update Technique</td>
<td>2233</td>
</tr>
<tr>
<td>66.4.2. Partial Match Algorithm</td>
<td>2234</td>
</tr>
<tr>
<td>66.5. GIN Tips and Tricks</td>
<td>2234</td>
</tr>
<tr>
<td>66.6. Limitations</td>
<td>2235</td>
</tr>
<tr>
<td>66.7. Examples</td>
<td>2235</td>
</tr>
<tr>
<td>67. BRIN Indexes</td>
<td>2236</td>
</tr>
<tr>
<td>67.1. Introduction</td>
<td>2236</td>
</tr>
<tr>
<td>67.1.1. Index Maintenance</td>
<td>2236</td>
</tr>
<tr>
<td>67.2. Built-in Operator Classes</td>
<td>2237</td>
</tr>
<tr>
<td>67.3. Extensibility</td>
<td>2238</td>
</tr>
<tr>
<td>68. Database Physical Storage</td>
<td>2241</td>
</tr>
<tr>
<td>68.1. Database File Layout</td>
<td>2241</td>
</tr>
<tr>
<td>68.2. TOAST</td>
<td>2243</td>
</tr>
<tr>
<td>68.2.1. Out-of-line, on-disk TOAST storage</td>
<td>2244</td>
</tr>
<tr>
<td>68.2.2. Out-of-line, in-memory TOAST storage</td>
<td>2245</td>
</tr>
<tr>
<td>68.3. Free Space Map</td>
<td>2246</td>
</tr>
<tr>
<td>68.4. Visibility Map</td>
<td>2246</td>
</tr>
</tbody>
</table>
Chapter 51. Overview of PostgreSQL Internals

Author

This chapter originated as part of [sim98], Stefan Simkovics' Master's Thesis prepared at Vienna University of Technology under the direction of O.Univ.Prof.Dr. Georg Gottlob and Univ.Ass. Mag. Katrin Seyr.

This chapter gives an overview of the internal structure of the backend of PostgreSQL. After having read the following sections you should have an idea of how a query is processed. This chapter does not aim to provide a detailed description of the internal operation of PostgreSQL, as such a document would be very extensive. Rather, this chapter is intended to help the reader understand the general sequence of operations that occur within the backend from the point at which a query is received, to the point at which the results are returned to the client.

51.1. The Path of a Query

Here we give a short overview of the stages a query has to pass in order to obtain a result.

1. A connection from an application program to the PostgreSQL server has to be established. The application program transmits a query to the server and waits to receive the results sent back by the server.

2. The parser stage checks the query transmitted by the application program for correct syntax and creates a query tree.

3. The rewrite system takes the query tree created by the parser stage and looks for any rules (stored in the system catalogs) to apply to the query tree. It performs the transformations given in the rule bodies.

One application of the rewrite system is in the realization of views. Whenever a query against a view (i.e., a virtual table) is made, the rewrite system rewrites the user's query to a query that accesses the base tables given in the view definition instead.

4. The planner/optimizer takes the (rewritten) query tree and creates a query plan that will be the input to the executor.

It does so by first creating all possible paths leading to the same result. For example if there is an index on a relation to be scanned, there are two paths for the scan. One possibility is a simple sequential scan and the other possibility is to use the index. Next the cost for the execution of each path is estimated and the cheapest path is chosen. The cheapest path is expanded into a complete plan that the executor can use.

5. The executor recursively steps through the plan tree and retrieves rows in the way represented by the plan. The executor makes use of the storage system while scanning relations, performs sorts and joins, evaluates qualifications and finally hands back the rows derived.

In the following sections we will cover each of the above listed items in more detail to give a better understanding of PostgreSQL's internal control and data structures.

51.2. How Connections are Established

PostgreSQL is implemented using a simple “process per user” client/server model. In this model there is one client process connected to exactly one server process. As we do not know ahead of time
how many connections will be made, we have to use a master process that spawns a new server process every time a connection is requested. This master process is called postgres and listens at a specified TCP/IP port for incoming connections. Whenever a request for a connection is detected the postgres process spawns a new server process. The server tasks communicate with each other using semaphores and shared memory to ensure data integrity throughout concurrent data access.

The client process can be any program that understands the PostgreSQL protocol described in Chapter 53. Many clients are based on the C-language library libpq, but several independent implementations of the protocol exist, such as the Java JDBC driver. Once a connection is established the client process can send a query to the backend (server). The query is transmitted using plain text, i.e., there is no parsing done in the frontend (client). The server parses the query, creates an execution plan, executes the plan and returns the retrieved rows to the client by transmitting them over the established connection.

51.3. The Parser Stage

The parser stage consists of two parts:

- The parser defined in gram.y and scan.l is built using the Unix tools bison and flex.
- The transformation process does modifications and augmentations to the data structures returned by the parser.

51.3.1. Parser

The parser has to check the query string (which arrives as plain text) for valid syntax. If the syntax is correct a parse tree is built up and handed back; otherwise an error is returned. The parser and lexer are implemented using the well-known Unix tools bison and flex.

The lexer is defined in the file scan.l and is responsible for recognizing identifiers, the SQL key words etc. For every key word or identifier that is found, a token is generated and handed to the parser.

The parser is defined in the file gram.y and consists of a set of grammar rules and actions that are executed whenever a rule is fired. The code of the actions (which is actually C code) is used to build up the parse tree.

The file scan.l is transformed to the C source file scan.c using the program flex and gram.y is transformed to gram.c using bison. After these transformations have taken place a normal C compiler can be used to create the parser. Never make any changes to the generated C files as they will be overwritten the next time flex or bison is called.

A detailed description of bison or the grammar rules given in gram.y would be beyond the scope of this paper. There are many books and documents dealing with flex and bison. You should be familiar with bison before you start to study the grammar given in gram.y otherwise you won’t understand what happens there.

51.3.2. Transformation Process

The parser stage creates a parse tree using only fixed rules about the syntactic structure of SQL. It does not make any lookups in the system catalogs, so there is no possibility to understand the detailed
semantics of the requested operations. After the parser completes, the transformation process takes the tree handed back by the parser as input and does the semantic interpretation needed to understand which tables, functions, and operators are referenced by the query. The data structure that is built to represent this information is called the query tree.

The reason for separating raw parsing from semantic analysis is that system catalog lookups can only be done within a transaction, and we do not wish to start a transaction immediately upon receiving a query string. The raw parsing stage is sufficient to identify the transaction control commands (BEGIN, ROLLBACK, etc), and these can then be correctly executed without any further analysis. Once we know that we are dealing with an actual query (such as SELECT or UPDATE), it is okay to start a transaction if we're not already in one. Only then can the transformation process be invoked.

The query tree created by the transformation process is structurally similar to the raw parse tree in most places, but it has many differences in detail. For example, a FuncCall node in the parse tree represents something that looks syntactically like a function call. This might be transformed to either a FuncExpr or Aggref node depending on whether the referenced name turns out to be an ordinary function or an aggregate function. Also, information about the actual data types of columns and expression results is added to the query tree.

51.4. The PostgreSQL Rule System

PostgreSQL supports a powerful rule system for the specification of views and ambiguous view updates. Originally the PostgreSQL rule system consisted of two implementations:

- The first one worked using row level processing and was implemented deep in the executor. The rule system was called whenever an individual row had been accessed. This implementation was removed in 1995 when the last official release of the Berkeley Postgres project was transformed into Postgres95.

- The second implementation of the rule system is a technique called query rewriting. The rewrite system is a module that exists between the parser stage and the planner/optimizer. This technique is still implemented.

The query rewriter is discussed in some detail in Chapter 41, so there is no need to cover it here. We will only point out that both the input and the output of the rewriter are query trees, that is, there is no change in the representation or level of semantic detail in the trees. Rewriting can be thought of as a form of macro expansion.

51.5. Planner/Optimizer

The task of the planner/optimizer is to create an optimal execution plan. A given SQL query (and hence, a query tree) can be actually executed in a wide variety of different ways, each of which will produce the same set of results. If it is computationally feasible, the query optimizer will examine each of these possible execution plans, ultimately selecting the execution plan that is expected to run the fastest.

Note

In some situations, examining each possible way in which a query can be executed would take an excessive amount of time and memory space. In particular, this occurs when executing queries involving large numbers of join operations. In order to determine a reasonable (not necessarily optimal) query plan in a reasonable amount of time, PostgreSQL uses a Genetic Query Optimizer (see Chapter 60) when the number of joins exceeds a threshold (see geqo_threshold).
The planner's search procedure actually works with data structures called *paths*, which are simply cutdown representations of plans containing only as much information as the planner needs to make its decisions. After the cheapest path is determined, a full-fledged *plan tree* is built to pass to the executor. This represents the desired execution plan in sufficient detail for the executor to run it. In the rest of this section we'll ignore the distinction between paths and plans.

### 51.5.1. Generating Possible Plans

The planner/optimizer starts by generating plans for scanning each individual relation (table) used in the query. The possible plans are determined by the available indexes on each relation. There is always the possibility of performing a sequential scan on a relation, so a sequential scan plan is always created. Assume an index is defined on a relation (for example a B-tree index) and a query contains the restriction `relation.attribute OPR constant`. If `relation.attribute` happens to match the key of the B-tree index and `OPR` is one of the operators listed in the index’s `operator class`, another plan is created using the B-tree index to scan the relation. If there are further indexes present and the restrictions in the query happen to match a key of an index, further plans will be considered. Index scan plans are also generated for indexes that have a sort ordering that can match the query’s `ORDER BY` clause (if any), or a sort ordering that might be useful for merge joining (see below).

If the query requires joining two or more relations, plans for joining relations are considered after all feasible plans have been found for scanning single relations. The three available join strategies are:

- **nested loop join**: The right relation is scanned once for every row found in the left relation. This strategy is easy to implement but can be very time consuming. (However, if the right relation can be scanned with an index scan, this can be a good strategy. It is possible to use values from the current row of the left relation as keys for the index scan of the right.)

- **merge join**: Each relation is sorted on the join attributes before the join starts. Then the two relations are scanned in parallel, and matching rows are combined to form join rows. This kind of join is more attractive because each relation has to be scanned only once. The required sorting might be achieved either by an explicit sort step, or by scanning the relation in the proper order using an index on the join key.

- **hash join**: The right relation is first scanned and loaded into a hash table, using its join attributes as hash keys. Next the left relation is scanned and the appropriate values of every row found are used as hash keys to locate the matching rows in the table.

When the query involves more than two relations, the final result must be built up by a tree of join steps, each with two inputs. The planner examines different possible join sequences to find the cheapest one.

If the query uses fewer than `geqo_threshold` relations, a near-exhaustive search is conducted to find the best join sequence. The planner preferentially considers joins between any two relations for which there exist a corresponding join clause in the `WHERE` qualification (i.e., for which a restriction like `where rel1.attr1=rel2.attr2` exists). Join pairs with no join clause are considered only when there is no other choice, that is, a particular relation has no available join clauses to any other relation. All possible plans are generated for every join pair considered by the planner, and the one that is (estimated to be) the cheapest is chosen.

When `geqo_threshold` is exceeded, the join sequences considered are determined by heuristics, as described in Chapter 60. Otherwise the process is the same.

The finished plan tree consists of sequential or index scans of the base relations, plus nested-loop, merge, or hash join nodes as needed, plus any auxiliary steps needed, such as sort nodes or aggregate-function calculation nodes. Most of these plan node types have the additional ability to do *selection* (discarding rows that do not meet a specified Boolean condition) and *projection* (computation of a derived column set based on given column values, that is, evaluation of scalar expressions where needed). One of the responsibilities of the planner is to attach selection conditions from the `WHERE` clause and computation of required output expressions to the most appropriate nodes of the plan tree.
51.6. Executor

The executor takes the plan created by the planner/optimizer and recursively processes it to extract the required set of rows. This is essentially a demand-pull pipeline mechanism. Each time a plan node is called, it must deliver one more row, or report that it is done delivering rows.

To provide a concrete example, assume that the top node is a MergeJoin node. Before any merge can be done two rows have to be fetched (one from each subplan). So the executor recursively calls itself to process the subplans (it starts with the subplan attached to lefttree). The new top node (the top node of the left subplan) is, let's say, a Sort node and again recursion is needed to obtain an input row. The child node of the Sort might be a SeqScan node, representing actual reading of a table. Execution of this node causes the executor to fetch a row from the table and return it up to the calling node. The Sort node will repeatedly call its child to obtain all the rows to be sorted. When the input is exhausted (as indicated by the child node returning a NULL instead of a row), the Sort code performs the sort, and finally is able to return its first output row, namely the first one in sorted order. It keeps the remaining rows stored so that it can deliver them in sorted order in response to later demands.

The MergeJoin node similarly demands the first row from its right subplan. Then it compares the two rows to see if they can be joined; if so, it returns a join row to its caller. On the next call, or immediately if it cannot join the current pair of inputs, it advances to the next row of one table or the other (depending on how the comparison came out), and again checks for a match. Eventually, one subplan or the other is exhausted, and the MergeJoin node returns NULL to indicate that no more join rows can be formed.

Complex queries can involve many levels of plan nodes, but the general approach is the same: each node computes and returns its next output row each time it is called. Each node is also responsible for applying any selection or projection expressions that were assigned to it by the planner.

The executor mechanism is used to evaluate all four basic SQL query types: SELECT, INSERT, UPDATE, and DELETE. For SELECT, the top-level executor code only needs to send each row returned by the query plan tree off to the client. For INSERT, each returned row is inserted into the target table specified for the INSERT. This is done in a special top-level plan node called ModifyTable. (A simple INSERT ... VALUES command creates a trivial plan tree consisting of a single Result node, which computes just one result row, and ModifyTable above it to perform the insertion. But INSERT ... SELECT can demand the full power of the executor mechanism.) For UPDATE, the planner arranges that each computed row includes all the updated column values, plus the TID (tuple ID, or row ID) of the original target row; this data is fed into a ModifyTable node, which uses the information to create a new updated row and mark the old row deleted. For DELETE, the only column that is actually returned by the plan is the TID, and the ModifyTable node simply uses the TID to visit each target row and mark it deleted.
Chapter 52. System Catalogs

The system catalogs are the place where a relational database management system stores schema metadata, such as information about tables and columns, and internal bookkeeping information. PostgreSQL's system catalogs are regular tables. You can drop and recreate the tables, add columns, insert and update values, and severely mess up your system that way. Normally, one should not change the system catalogs by hand, there are normally SQL commands to do that. (For example, CREATE DATABASE inserts a row into the pg_database catalog — and actually creates the database on disk.) There are some exceptions for particularly esoteric operations, but many of those have been made available as SQL commands over time, and so the need for direct manipulation of the system catalogs is ever decreasing.

52.1. Overview

Table 52.1 lists the system catalogs. More detailed documentation of each catalog follows below.

Most system catalogs are copied from the template database during database creation and are thereafter database-specific. A few catalogs are physically shared across all databases in a cluster; these are noted in the descriptions of the individual catalogs.

<table>
<thead>
<tr>
<th>Catalog Name</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_aggregate</td>
<td>aggregate functions</td>
</tr>
<tr>
<td>pg_am</td>
<td>index access methods</td>
</tr>
<tr>
<td>pg_amop</td>
<td>access method operators</td>
</tr>
<tr>
<td>pg_amproc</td>
<td>access method support functions</td>
</tr>
<tr>
<td>pg_attrdef</td>
<td>column default values</td>
</tr>
<tr>
<td>pg_attribute</td>
<td>table columns (&quot;attributes&quot;)</td>
</tr>
<tr>
<td>pg_authid</td>
<td>authorization identifiers (roles)</td>
</tr>
<tr>
<td>pg_auth_members</td>
<td>authorization identifier membership relationships</td>
</tr>
<tr>
<td>pg_cast</td>
<td>casts (data type conversions)</td>
</tr>
<tr>
<td>pg_class</td>
<td>tables, indexes, sequences, views (&quot;relations&quot;)</td>
</tr>
<tr>
<td>pg_collation</td>
<td>collations (locale information)</td>
</tr>
<tr>
<td>pg_constraint</td>
<td>check constraints, unique constraints, primary key constraints, foreign key constraints</td>
</tr>
<tr>
<td>pg_conversion</td>
<td>encoding conversion information</td>
</tr>
<tr>
<td>pg_database</td>
<td>databases within this database cluster</td>
</tr>
<tr>
<td>pg_db_role_setting</td>
<td>per-role and per-database settings</td>
</tr>
<tr>
<td>pg_default_acl</td>
<td>default privileges for object types</td>
</tr>
<tr>
<td>pg_depend</td>
<td>dependencies between database objects</td>
</tr>
<tr>
<td>pg_description</td>
<td>descriptions or comments on database objects</td>
</tr>
<tr>
<td>pg_enum</td>
<td>enum label and value definitions</td>
</tr>
<tr>
<td>pg_event_trigger</td>
<td>event triggers</td>
</tr>
<tr>
<td>pg_extension</td>
<td>installed extensions</td>
</tr>
<tr>
<td>pg_foreign_data_wrapper</td>
<td>foreign-data wrapper definitions</td>
</tr>
<tr>
<td>pg_foreign_server</td>
<td>foreign server definitions</td>
</tr>
<tr>
<td>pg_foreign_table</td>
<td>additional foreign table information</td>
</tr>
<tr>
<td>Catalog Name</td>
<td>Purpose</td>
</tr>
<tr>
<td>------------------------------</td>
<td>--------------------------------------------</td>
</tr>
<tr>
<td>pg_index</td>
<td>additional index information</td>
</tr>
<tr>
<td>pg_inherits</td>
<td>table inheritance hierarchy</td>
</tr>
<tr>
<td>pg_init_privs</td>
<td>object initial privileges</td>
</tr>
<tr>
<td>pg_language</td>
<td>languages for writing functions</td>
</tr>
<tr>
<td>pg_largeobject</td>
<td>data pages for large objects</td>
</tr>
<tr>
<td>pg_largeobject_metadata</td>
<td>metadata for large objects</td>
</tr>
<tr>
<td>pg_namespace</td>
<td>schemas</td>
</tr>
<tr>
<td>pg_opclass</td>
<td>access method operator classes</td>
</tr>
<tr>
<td>pg_operator</td>
<td>operators</td>
</tr>
<tr>
<td>pg_opfamily</td>
<td>access method operator families</td>
</tr>
<tr>
<td>pg_partitioned_table</td>
<td>information about partition key of tables</td>
</tr>
<tr>
<td>pg_pltemplate</td>
<td>template data for procedural languages</td>
</tr>
<tr>
<td>pg_policy</td>
<td>row-security policies</td>
</tr>
<tr>
<td>pg_proc</td>
<td>functions and procedures</td>
</tr>
<tr>
<td>pg_publication</td>
<td>publications for logical replication</td>
</tr>
<tr>
<td>pg_publication_rel</td>
<td>relation to publication mapping</td>
</tr>
<tr>
<td>pg_range</td>
<td>information about range types</td>
</tr>
<tr>
<td>pg_replication_origin</td>
<td>registered replication origins</td>
</tr>
<tr>
<td>pg_rewrite</td>
<td>query rewrite rules</td>
</tr>
<tr>
<td>pg_seclabel</td>
<td>security labels on database objects</td>
</tr>
<tr>
<td>pg_sequence</td>
<td>information about sequences</td>
</tr>
<tr>
<td>pg_shdepend</td>
<td>dependencies on shared objects</td>
</tr>
<tr>
<td>pg_shdescription</td>
<td>comments on shared objects</td>
</tr>
<tr>
<td>pg_shseclabel</td>
<td>security labels on shared database objects</td>
</tr>
<tr>
<td>pg_statistic</td>
<td>planner statistics</td>
</tr>
<tr>
<td>pg_statistic_ext</td>
<td>extended planner statistics</td>
</tr>
<tr>
<td>pg_subscriptions</td>
<td>logical replication subscriptions</td>
</tr>
<tr>
<td>pg_subscription_rel</td>
<td>relation state for subscriptions</td>
</tr>
<tr>
<td>pg_tablespace</td>
<td>tablespaces within this database cluster</td>
</tr>
<tr>
<td>pg_transform</td>
<td>transforms (data type to procedural language conversions)</td>
</tr>
<tr>
<td>pg_trigger</td>
<td>triggers</td>
</tr>
<tr>
<td>pg_ts_config</td>
<td>text search configurations</td>
</tr>
<tr>
<td>pg_ts_config_map</td>
<td>text search configurations' token mappings</td>
</tr>
<tr>
<td>pg_ts_dict</td>
<td>text search dictionaries</td>
</tr>
<tr>
<td>pg_ts_parser</td>
<td>text search parsers</td>
</tr>
<tr>
<td>pg_ts_template</td>
<td>text search templates</td>
</tr>
<tr>
<td>pg_type</td>
<td>data types</td>
</tr>
<tr>
<td>pg_user_mapping</td>
<td>mappings of users to foreign servers</td>
</tr>
</tbody>
</table>

52.2. pg_aggregate
The catalog `pg_aggregate` stores information about aggregate functions. An aggregate function is a function that operates on a set of values (typically one column from each row that matches a query condition) and returns a single value computed from all these values. Typical aggregate functions are `sum`, `count`, and `max`. Each entry in `pg_aggregate` is an extension of an entry in `pg_proc`. The `pg_proc` entry carries the aggregate's name, input and output data types, and other information that is similar to ordinary functions.

### Table 52.2. `pg_aggregate` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>aggfnoid</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>pg_proc OID of the aggregate function</td>
</tr>
<tr>
<td>aggkind</td>
<td>char</td>
<td></td>
<td>Aggregate kind: n for “normal” aggregates, o for “ordered-set” aggregates, or h for “hypothetical-set” aggregates</td>
</tr>
<tr>
<td>aggnumdirectargs</td>
<td>int2</td>
<td></td>
<td>Number of direct (non-aggregated) arguments of an ordered-set or hypothetical-set aggregate, counting a variadic array as one argument. If equal to <code>pronargs</code>, the aggregate must be variadic and the variadic array describes the aggregated arguments as well as the final direct arguments. Always zero for normal aggregates.</td>
</tr>
<tr>
<td>aggtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Transition function</td>
</tr>
<tr>
<td>aggfinafinaln</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function (zero if none)</td>
</tr>
<tr>
<td>aggcombinefn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Combine function (zero if none)</td>
</tr>
<tr>
<td>aggserializefn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Serialization function (zero if none)</td>
</tr>
<tr>
<td>aggserializefn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Deserialization function (zero if none)</td>
</tr>
<tr>
<td>aggfinaln</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggfinaln</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Forward transition function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Inverse transition function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmfinaln</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>agffinalextra</td>
<td>bool</td>
<td></td>
<td>True to pass extra dummy arguments to aggfinaln</td>
</tr>
</tbody>
</table>

2001
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>aggmfinalextra</td>
<td>bool</td>
<td></td>
<td>True to pass extra dummy arguments to aggmfinalfn</td>
</tr>
<tr>
<td>aggfinalmodify</td>
<td>char</td>
<td></td>
<td>Whether aggfinalfn modifies the transition state value: r if it is read-only, s if the aggransfn cannot be applied after the aggfinalfn, or w if it writes on the value</td>
</tr>
<tr>
<td>aggmfinalmodify</td>
<td>char</td>
<td></td>
<td>Like aggfinalmodify, but for the aggmfinalfn</td>
</tr>
<tr>
<td>aggsortop</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Associated sort operator (zero if none)</td>
</tr>
<tr>
<td>aggtranstype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the aggregate function’s internal transition (state) data</td>
</tr>
<tr>
<td>aggransspace</td>
<td>int4</td>
<td></td>
<td>Approximate average size (in bytes) of the transition state data, or zero to use a default estimate</td>
</tr>
<tr>
<td>aggmtranstype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the aggregate function’s internal transition (state) data for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmtransspace</td>
<td>int4</td>
<td></td>
<td>Approximate average size (in bytes) of the transition state data for moving-aggregate mode, or zero to use a default estimate</td>
</tr>
<tr>
<td>agginitialval</td>
<td>text</td>
<td></td>
<td>The initial value of the transition state. This is a text field containing the initial value in its external string representation. If this field is null, the transition state value starts out null.</td>
</tr>
<tr>
<td>aggmminimalval</td>
<td>text</td>
<td></td>
<td>The initial value of the transition state for moving-aggregate mode. This is a text field containing the initial value in its external string representation. If this field is null, the transition state value starts out null.</td>
</tr>
</tbody>
</table>
New aggregate functions are registered with the CREATE AGGREGATE command. See Section 38.11 for more information about writing aggregate functions and the meaning of the transition functions, etc.

52.3. pg_am

The catalog pg_am stores information about relation access methods. There is one row for each access method supported by the system. Currently, only indexes have access methods. The requirements for index access methods are discussed in detail in Chapter 61.

Table 52.3. pg_am Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>amname</td>
<td>name</td>
<td></td>
<td>Name of the access method</td>
</tr>
<tr>
<td>amhandler</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of a handler function that is responsible for supplying information about the access method</td>
</tr>
<tr>
<td>amtype</td>
<td>char</td>
<td></td>
<td>Currently always i to indicate an index access method; other values may be allowed in future</td>
</tr>
</tbody>
</table>

Note

Before PostgreSQL 9.6, pg_am contained many additional columns representing properties of index access methods. That data is now only directly visible at the C code level. However, pg_index_column_has_property() and related functions have been added to allow SQL queries to inspect index access method properties; see Table 9.63.

52.4. pg_amop

The catalog pg_amop stores information about operators associated with access method operator families. There is one row for each operator that is a member of an operator family. A family member can be either a search operator or an ordering operator. An operator can appear in more than one family, but cannot appear in more than one search position nor more than one ordering position within a family. (It is allowed, though unlikely, for an operator to be used for both search and ordering purposes.)

Table 52.4. pg_amop Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>amopfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The operator family this entry is for</td>
</tr>
<tr>
<td>amoplefttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Left-hand input data type of operator</td>
</tr>
<tr>
<td>amoprighttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Right-hand input data type of operator</td>
</tr>
<tr>
<td>amopstrategy</td>
<td>int2</td>
<td></td>
<td>Operator strategy number</td>
</tr>
<tr>
<td>amoppurpose</td>
<td>char</td>
<td></td>
<td>Operator purpose, either s for search or o for ordering</td>
</tr>
<tr>
<td>amopopr</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>OID of the operator</td>
</tr>
<tr>
<td>amopmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator family is for</td>
</tr>
<tr>
<td>amopsortfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The B-tree operator family this entry sorts according to, if an ordering operator; zero if a search operator</td>
</tr>
</tbody>
</table>

A “search” operator entry indicates that an index of this operator family can be searched to find all rows satisfying `WHERE indexed_column operator constant`. Obviously, such an operator must return `boolean`, and its left-hand input type must match the index’s column data type.

An “ordering” operator entry indicates that an index of this operator family can be scanned to return rows in the order represented by `ORDER BY indexed_column operator constant`. Such an operator could return any sortable data type, though again its left-hand input type must match the index’s column data type. The exact semantics of the `ORDER BY` are specified by the `amopsortfamily` column, which must reference a B-tree operator family for the operator’s result type.

**Note**

At present, it's assumed that the sort order for an ordering operator is the default for the referenced operator family, i.e., `ASC NULLS LAST`. This might someday be relaxed by adding additional columns to specify sort options explicitly.

An entry’s `amopmethod` must match the `opfmethod` of its containing operator family (including `amopmethod` here is an intentional denormalization of the catalog structure for performance reasons). Also, `amoplefttype` and `amoprighttype` must match the `oprleft` and `oppright` fields of the referenced `pg_operator` entry.

**52.5. pg_amproc**

The catalog `pg_amproc` stores information about support functions associated with access method operator families. There is one row for each support function belonging to an operator family.

**Table 52.5. pg_amproc Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>amprocfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The operator family this entry is for</td>
</tr>
<tr>
<td>amproclefttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Left-hand input data type of associated operator</td>
</tr>
<tr>
<td>amprocriggtype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Right-hand input data type of associated operator</td>
</tr>
<tr>
<td>amprocnum</td>
<td>int2</td>
<td></td>
<td>Support function number</td>
</tr>
<tr>
<td>amproc</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the function</td>
</tr>
</tbody>
</table>

The usual interpretation of the amproclefttype and amprocriggtype fields is that they identify the left and right input types of the operator(s) that a particular support function supports. For some access methods these match the input data type(s) of the support function itself, for others not. There is a notion of “default” support functions for an index, which are those with amproclefttype and amprocriggtype both equal to the index operator class's opcintype.

### 52.6. `pg_attrdef`

The catalog `pg_attrdef` stores column default values. The main information about columns is stored in `pg_attribute` (see below). Only columns that explicitly specify a default value (when the table is created or the column is added) will have an entry here.

#### Table 52.6. `pg_attrdef` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>adrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this column belongs to</td>
</tr>
<tr>
<td>adnum</td>
<td>int2</td>
<td>pg_attribute.atnum</td>
<td>The number of the column</td>
</tr>
<tr>
<td>adbin</td>
<td>pg_node_tree</td>
<td></td>
<td>The internal representation of the column default value</td>
</tr>
<tr>
<td>adsrcre</td>
<td>text</td>
<td></td>
<td>A human-readable representation of the default value</td>
</tr>
</tbody>
</table>

The adsrcre field is historical, and is best not used, because it does not track outside changes that might affect the representation of the default value. Reverse-compiling the adbin field (with `pg_get_expr` for example) is a better way to display the default value.

### 52.7. `pg_attribute`

The catalog `pg_attribute` stores information about table columns. There will be exactly one `pg_attribute` row for every column in every table in the database. (There will also be attribute entries for indexes, and indeed all objects that have `pg_class` entries.)

The term attribute is equivalent to column and is used for historical reasons.
### Table 52.7. pg_attribute Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>attrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this column belongs to</td>
</tr>
<tr>
<td>attname</td>
<td>name</td>
<td></td>
<td>The column name</td>
</tr>
<tr>
<td>atttypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The data type of this column</td>
</tr>
<tr>
<td>attstattraget</td>
<td>int4</td>
<td></td>
<td>attstattraget controls the level of detail of statistics accumulated for this column by ANALYZE. A zero value indicates that no statistics should be collected. A negative value says to use the system default statistics target. The exact meaning of positive values is data type-dependent. For scalar data types, attstattraget is both the target number of “most common values” to collect, and the target number of histogram bins to create.</td>
</tr>
<tr>
<td>attlen</td>
<td>int2</td>
<td></td>
<td>A copy of pg_type.typlen of this column’s type</td>
</tr>
<tr>
<td>attnum</td>
<td>int2</td>
<td></td>
<td>The number of the column. Ordinary columns are numbered from 1 up. System columns, such as oid, have (arbitrary) negative numbers.</td>
</tr>
<tr>
<td>attndims</td>
<td>int4</td>
<td></td>
<td>Number of dimensions, if the column is an array type; otherwise 0. (Presently, the number of dimensions of an array is not enforced, so any nonzero value effectively means “it’s an array”.)</td>
</tr>
<tr>
<td>attcacheoff</td>
<td>int4</td>
<td></td>
<td>Always -1 in storage, but when loaded into a row descriptor in memory this might be updated to cache the offset of the attribute within the row</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>--------</td>
<td>------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>atttypmod</td>
<td>int4</td>
<td></td>
<td>atttypmod records type-specific data supplied at table creation time (for example, the maximum length of a varchar column). It is passed to type-specific input functions and length coercion functions. The value will generally be -1 for types that do not need atttypmod.</td>
</tr>
<tr>
<td>attbyval</td>
<td>bool</td>
<td></td>
<td>A copy of pg_type.typbyval of this column’s type.</td>
</tr>
<tr>
<td>attstorage</td>
<td>char</td>
<td></td>
<td>Normally a copy of pg_type.typstorage of this column’s type. For TOAST-able data types, this can be altered after column creation to control storage policy.</td>
</tr>
<tr>
<td>attalign</td>
<td>char</td>
<td></td>
<td>A copy of pg_type.typalign of this column’s type.</td>
</tr>
<tr>
<td>attnotnull</td>
<td>bool</td>
<td></td>
<td>This represents a not-null constraint.</td>
</tr>
<tr>
<td>atthasdef</td>
<td>bool</td>
<td></td>
<td>This column has a default value, in which case there will be a corresponding entry in the pg_attrdef catalog that actually defines the value.</td>
</tr>
<tr>
<td>atthasmissing</td>
<td>bool</td>
<td></td>
<td>This column has a value which is used where the column is entirely missing from the row, as happens when a column is added with a non-volatile DEFAULT value after the row is created. The actual value used is stored in the atthasmissingval column.</td>
</tr>
<tr>
<td>attidentity</td>
<td>char</td>
<td></td>
<td>If a zero byte (&quot;'&quot;), then not an identity column. Otherwise, a = generated always, d = generated by default.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>--------------</td>
<td>------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>attisdropped</td>
<td>bool</td>
<td></td>
<td>This column has been dropped and is no longer valid. A dropped column is still physically present in the table, but is ignored by the parser and so cannot be accessed via SQL.</td>
</tr>
<tr>
<td>attislocal</td>
<td>bool</td>
<td></td>
<td>This column is defined locally in the relation. Note that a column can be locally defined and inherited simultaneously.</td>
</tr>
<tr>
<td>attinhcount</td>
<td>int4</td>
<td></td>
<td>The number of direct ancestors this column has. A column with a nonzero number of ancestors cannot be dropped nor renamed.</td>
</tr>
<tr>
<td>attcollation</td>
<td>oid</td>
<td>pg_collation.oid</td>
<td>The defined collation of the column, or zero if the column is not of a collatable data type.</td>
</tr>
<tr>
<td>attacl</td>
<td>aclitem[]</td>
<td></td>
<td>Column-level access privileges, if any have been granted specifically on this column.</td>
</tr>
<tr>
<td>attoptions</td>
<td>text[]</td>
<td></td>
<td>Attribute-level options, as “keyword=value” strings.</td>
</tr>
<tr>
<td>attfdwoptions</td>
<td>text[]</td>
<td></td>
<td>Attribute-level foreign data wrapper options, as “keyword=value” strings.</td>
</tr>
<tr>
<td>attmissingval</td>
<td>anyarray</td>
<td></td>
<td>This column has a one element array containing the value used when the column is entirely missing from the row, as happens when the column is added with a non-volatile DEFAULT value after the row is created. The value is only used when atthasmissing is true. If there is no value the column is null.</td>
</tr>
</tbody>
</table>

In a dropped column’s `pg_attribute` entry, `atttypid` is reset to zero, but `attlen` and the other fields copied from `pg_type` are still valid. This arrangement is needed to cope with the situation where the dropped column’s data type was later dropped, and so there is no `pg_type` row anymore. `attlen` and the other fields can be used to interpret the contents of a row of the table.
52.8. pg_authid

The catalog `pg_authid` contains information about database authorization identifiers (roles). A role subsumes the concepts of “users” and “groups”. A user is essentially just a role with the `rolcanlogin` flag set. Any role (with or without `rolcanlogin`) can have other roles as members; see `pg_auth_members`.

Since this catalog contains passwords, it must not be publicly readable. `pg_roles` is a publicly readable view on `pg_authid` that blanks out the password field.

Chapter 21 contains detailed information about user and privilege management.

Because user identities are cluster-wide, `pg_authid` is shared across all databases of a cluster: there is only one copy of `pg_authid` per cluster, not one per database.

**Table 52.8. pg_authid Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>rolname</td>
<td>name</td>
<td>Role name</td>
</tr>
<tr>
<td>rolsuper</td>
<td>bool</td>
<td>Role has superuser privileges</td>
</tr>
<tr>
<td>rolinherit</td>
<td>bool</td>
<td>Role automatically inherits privileges of roles it is a member of</td>
</tr>
<tr>
<td>rolcreaterole</td>
<td>bool</td>
<td>Role can create more roles</td>
</tr>
<tr>
<td>rolcreatedb</td>
<td>bool</td>
<td>Role can create databases</td>
</tr>
<tr>
<td>rolcanlogin</td>
<td>bool</td>
<td>Role can log in. That is, this role can be given as the initial session authorization identifier</td>
</tr>
<tr>
<td>rolreplication</td>
<td>bool</td>
<td>Role is a replication role. A replication role can initiate replication connections and create and drop replication slots.</td>
</tr>
<tr>
<td>rolbypassrls</td>
<td>bool</td>
<td>Role bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>rolconnlimit</td>
<td>int4</td>
<td>For roles that can log in, this sets maximum number of concurrent connections this role can make. -1 means no limit.</td>
</tr>
<tr>
<td>rolpassword</td>
<td>text</td>
<td>Password (possibly encrypted); null if none. The format depends on the form of encryption used.</td>
</tr>
<tr>
<td>rolvaliduntil</td>
<td>timestamptz</td>
<td>Password expiry time (only used for password authentication); null if no expiration</td>
</tr>
</tbody>
</table>

For an MD5 encrypted password, the `rolpassword` column will begin with the string `md5` followed by a 32-character hexadecimal MD5 hash. The MD5 hash will be of the user’s password concatenated to their user name. For example, if user `joe` has password `xyzzy`, PostgreSQL will store the MD5 hash of `xyzzyjoe`.

If the password is encrypted with SCRAM-SHA-256, it has the format:
System Catalogs

SCRAM-SHA-256$<iteration count>:<salt>:<StoredKey>:<ServerKey>

where salt, StoredKey and ServerKey are in Base64 encoded format. This format is the same as that specified by RFC 5803.

A password that does not follow either of those formats is assumed to be unencrypted.

52.9. pg_auth_members

The catalog pg_auth_members shows the membership relations between roles. Any non-circular set of relationships is allowed.

Because user identities are cluster-wide, pg_auth_members is shared across all databases of a cluster: there is only one copy of pg_auth_members per cluster, not one per database.

Table 52.9. pg_auth_members Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>roleid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of a role that has a member</td>
</tr>
<tr>
<td>member</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of a role that is a member of roleid</td>
</tr>
<tr>
<td>grantor</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of the role that granted this membership</td>
</tr>
<tr>
<td>admin_option</td>
<td>bool</td>
<td></td>
<td>True if member can grant membership in roleid to others</td>
</tr>
</tbody>
</table>

52.10. pg_cast

The catalog pg_cast stores data type conversion paths, both built-in and user-defined.

It should be noted that pg_cast does not represent every type conversion that the system knows how to perform; only those that cannot be deduced from some generic rule. For example, casting between a domain and its base type is not explicitly represented in pg_cast. Another important exception is that “automatic I/O conversion casts”, those performed using a data type's own I/O functions to convert to or from text or other string types, are not explicitly represented in pg_cast.

Table 52.10. pg_cast Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>castsource</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the source data type</td>
</tr>
<tr>
<td>casttarget</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the target data type</td>
</tr>
<tr>
<td>castfunc</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use to perform this cast. Zero is stored if the cast method doesn't require a function.</td>
</tr>
<tr>
<td>castcontext</td>
<td>char</td>
<td></td>
<td>Indicates what contexts the cast can be invoked</td>
</tr>
</tbody>
</table>
The cast functions listed in `pg_cast` must always take the cast source type as their first argument type, and return the cast destination type as their result type. A cast function can have up to three arguments. The second argument, if present, must be type `integer`; it receives the type modifier associated with the destination type, or -1 if there is none. The third argument, if present, must be type `boolean`; it receives `true` if the cast is an explicit cast, `false` otherwise.

It is legitimate to create a `pg_cast` entry in which the source and target types are the same, if the associated function takes more than one argument. Such entries represent “length coercion functions” that coerce values of the type to be legal for a particular type modifier value.

When a `pg_cast` entry has different source and target types and a function that takes more than one argument, it represents converting from one type to another and applying a length coercion in a single step. When no such entry is available, coercion to a type that uses a type modifier involves two steps, one to convert between data types and a second to apply the modifier.

## 52.11. `pg_class`

The catalog `pg_class` catalogs tables and most everything else that has columns or is otherwise similar to a table. This includes indexes (but see also `pg_index`), sequences (but see also `pg_sequence`), views, materialized views, composite types, and TOAST tables; see `relkind`. Below, when we mean all of these kinds of objects we speak of “relations”. Not all columns are meaningful for all relation types.

### Table 52.11. `pg_class` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>relnamespac</td>
<td>oid</td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace that contains this relation</td>
</tr>
<tr>
<td>relnamenspac</td>
<td>name</td>
<td></td>
<td>Name of the table, index, view, etc.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>----------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>reltype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The OID of the data type that corresponds to this table’s row type, if any (zero for indexes, which have no pg_type entry)</td>
</tr>
<tr>
<td>reloftype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>For typed tables, the OID of the underlying composite type, zero for all other relations</td>
</tr>
<tr>
<td>relowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the relation</td>
</tr>
<tr>
<td>relam</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>If this is an index, the access method used (B-tree, hash, etc.)</td>
</tr>
<tr>
<td>relfilenode</td>
<td>oid</td>
<td></td>
<td>Name of the on-disk file of this relation; zero means this is a “mapped” relation whose disk file name is determined by low-level state</td>
</tr>
<tr>
<td>reltablespace</td>
<td>oid</td>
<td>pg_table-space.oid</td>
<td>The tablespace in which this relation is stored. If zero, the database's default tablespace is implied. (Not meaningful if the relation has no on-disk file.)</td>
</tr>
<tr>
<td>relpages</td>
<td>int4</td>
<td></td>
<td>Size of the on-disk representation of this table in pages (of size BLCKSZ). This is only an estimate used by the planner. It is updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX.</td>
</tr>
<tr>
<td>reltuples</td>
<td>float4</td>
<td></td>
<td>Number of live rows in the table. This is only an estimate used by the planner. It is updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX.</td>
</tr>
<tr>
<td>relallvisible</td>
<td>int4</td>
<td></td>
<td>Number of pages that are marked all-visible in the table's visibility map. This is only an estimate used by the planner. It is updated by VACUUM,</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>-------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>reltoastrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the TOAST table associated with this table, 0 if none. The TOAST table stores large attributes “out of line” in a secondary table.</td>
</tr>
<tr>
<td>relhasindex</td>
<td>bool</td>
<td></td>
<td>True if this is a table and it has (or recently had) any indexes</td>
</tr>
<tr>
<td>relisshared</td>
<td>bool</td>
<td></td>
<td>True if this table is shared across all databases in the cluster. Only certain system catalogs (such as pg_database) are shared.</td>
</tr>
<tr>
<td>relpersistence</td>
<td>char</td>
<td></td>
<td>p = permanent table, u = unlogged table, t = temporary table</td>
</tr>
<tr>
<td>relkind</td>
<td>char</td>
<td></td>
<td>r = ordinary table, i = index, S = sequence, t = TOAST table, v = view, m = materialized view, c = composite type, f = foreign table, p = partitioned table, I = partitioned index</td>
</tr>
<tr>
<td>relnatts</td>
<td>int2</td>
<td></td>
<td>Number of user columns in the relation (system columns not counted). There must be this many corresponding entries in pg_attribute. See also pg_attribute.attnum.</td>
</tr>
<tr>
<td>relchecks</td>
<td>int2</td>
<td></td>
<td>Number of CHECK constraints on the table; see pg_constraint catalog</td>
</tr>
<tr>
<td>relhasoids</td>
<td>bool</td>
<td></td>
<td>True if we generate an OID for each row of the relation</td>
</tr>
<tr>
<td>relhasrules</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) rules; see pg_rewrite catalog</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>-----------</td>
<td>-------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>relhastriggers</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) triggers; see pg_trigger catalog</td>
</tr>
<tr>
<td>relhassubclass</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) any inheritance children</td>
</tr>
<tr>
<td>relrowsecurity</td>
<td>bool</td>
<td></td>
<td>True if table has row level security enabled; see pg_policy catalog</td>
</tr>
<tr>
<td>relforcerowsecurity</td>
<td>bool</td>
<td></td>
<td>True if row level security (when enabled) will also apply to table owner; see pg_policy catalog</td>
</tr>
<tr>
<td>relispopulated</td>
<td>bool</td>
<td></td>
<td>True if relation is populated (this is true for all relations other than some materialized views)</td>
</tr>
<tr>
<td>relreplident</td>
<td>char</td>
<td></td>
<td>Columns used to form “replica identity” for rows: ( d = \text{default (primary key, if any)} ), ( n = \text{nothing} ), ( f = \text{all columns} ), ( i = \text{index with indisreplident set, or default} )</td>
</tr>
<tr>
<td>relispartition</td>
<td>bool</td>
<td></td>
<td>True if table or index is a partition</td>
</tr>
<tr>
<td>relrewrite</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>For new relations being written during a DDL operation that requires a table rewrite, this contains the OID of the original relation; otherwise 0. That state is only visible internally; this field should never contain anything other than 0 for a user-visible relation.</td>
</tr>
</tbody>
</table>
| relfrozenxid       | xid       |                   | All transaction IDs before this one have been replaced with a permanent (“frozen”) transaction ID in this table. This is used to track whether the table needs to be vacuumed in order to prevent transaction ID wraparound or to allow pg_x-
Several of the Boolean flags in `pg_class` are maintained lazily: they are guaranteed to be true if that's the correct state, but may not be reset to false immediately when the condition is no longer true. For example, `relhasindex` is set by `CREATE INDEX`, but it is never cleared by `DROP INDEX`. Instead, `VACUUM` clears `relhasindex` if it finds the table has no indexes. This arrangement avoids race conditions and improves concurrency.

### 52.12. `pg_collation`

The catalog `pg_collation` describes the available collations, which are essentially mappings from an SQL name to operating system locale categories. See Section 23.2 for more information.

<table>
<thead>
<tr>
<th>Table 52.12. <code>pg_collation</code> Columns</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Name</strong></td>
</tr>
<tr>
<td><code>oid</code></td>
</tr>
<tr>
<td><code>collname</code></td>
</tr>
<tr>
<td><code>collnamespace</code></td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>collowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the collation</td>
</tr>
<tr>
<td>collprovider</td>
<td>char</td>
<td></td>
<td>Provider of the collation: d = database default, c = libc, i = icu</td>
</tr>
<tr>
<td>collencoding</td>
<td>int4</td>
<td></td>
<td>Encoding in which the collation is applicable, or -1 if it works for any encoding</td>
</tr>
<tr>
<td>collcollate</td>
<td>name</td>
<td></td>
<td>LC_COLLATE for this collation object</td>
</tr>
<tr>
<td>collctype</td>
<td>name</td>
<td></td>
<td>LC_CTYPE for this collation object</td>
</tr>
<tr>
<td>collversion</td>
<td>text</td>
<td></td>
<td>Provider-specific version of the collation. This is recorded when the collation is created and then checked when it is used, to detect changes in the collation definition that could lead to data corruption.</td>
</tr>
</tbody>
</table>

Note that the unique key on this catalog is (collname, collencoding, collnamespace) not just (collname, collnamespace). PostgreSQL generally ignores all collations that do not have collencoding equal to either the current database’s encoding or -1, and creation of new entries with the same name as an entry with collencoding = -1 is forbidden. Therefore it is sufficient to use a qualified SQL name (schema.name) to identify a collation, even though this is not unique according to the catalog definition. The reason for defining the catalog this way is that initdb fills it in at cluster initialization time with entries for all locales available on the system, so it must be able to hold entries for all encodings that might ever be used in the cluster.

In the template0 database, it could be useful to create collations whose encoding does not match the database encoding, since they could match the encodings of databases later cloned from template0. This would currently have to be done manually.

52.13. pg_constraint

The catalog pg_constraint stores check, primary key, unique, foreign key, and exclusion constraints on tables. (Column constraints are not treated specially. Every column constraint is equivalent to some table constraint.) Not-null constraints are represented in the pg_attribute catalog, not here.

User-defined constraint triggers (created with CREATE CONSTRAINT TRIGGER) also give rise to an entry in this table.

Check constraints on domains are stored here, too.

Table 52.13. pg_constraint Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>---------</td>
<td>---------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>conname</td>
<td>name</td>
<td></td>
<td>Constraint name (not necessarily unique!)</td>
</tr>
<tr>
<td>connamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this constraint</td>
</tr>
<tr>
<td>contype</td>
<td>char</td>
<td></td>
<td>c = check constraint, f = foreign key constraint, p = primary key constraint, u = unique constraint, t = constraint trigger, x = exclusion constraint</td>
</tr>
<tr>
<td>condeferrable</td>
<td>bool</td>
<td></td>
<td>Is the constraint deferrable?</td>
</tr>
<tr>
<td>condeferred</td>
<td>bool</td>
<td></td>
<td>Is the constraint deferred by default?</td>
</tr>
<tr>
<td>convalidated</td>
<td>bool</td>
<td></td>
<td>Has the constraint been validated? Currently, can only be false for foreign keys and CHECK constraints</td>
</tr>
<tr>
<td>conrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this constraint is on; 0 if not a table constraint</td>
</tr>
<tr>
<td>contypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The domain this constraint is on; 0 if not a domain constraint</td>
</tr>
<tr>
<td>conindid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The index supporting this constraint, if it’s a unique, primary key, foreign key, or exclusion constraint; else 0</td>
</tr>
<tr>
<td>conparentid</td>
<td>oid</td>
<td>pg_constraint.t.oid</td>
<td>The corresponding constraint in the parent partitioned table, if this is a constraint in a partition; else 0</td>
</tr>
<tr>
<td>confreloid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>If a foreign key, the referenced table; else 0</td>
</tr>
<tr>
<td>confupdtype</td>
<td>char</td>
<td></td>
<td>Foreign key update action code: a = no action, r = restrict, c = cascade, n = set null, d = set default</td>
</tr>
<tr>
<td>confdeltype</td>
<td>char</td>
<td></td>
<td>Foreign key deletion action code: a = no action, r = restrict, c = cascade, n = set null, d = set default</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>---------------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>confmatchtype</td>
<td>char</td>
<td></td>
<td>Foreign key match type: f = full, p = partial, s = simple</td>
</tr>
<tr>
<td>conislocal</td>
<td>bool</td>
<td></td>
<td>This constraint is defined locally for the relation. Note that a constraint can be locally defined and inherited simultaneously.</td>
</tr>
<tr>
<td>coninhcount</td>
<td>int4</td>
<td></td>
<td>The number of direct inheritance ancestors this constraint has. A constraint with a non-zero number of ancestors cannot be dropped nor renamed.</td>
</tr>
<tr>
<td>connoinherit</td>
<td>bool</td>
<td></td>
<td>This constraint is defined locally for the relation. It is a non-inheritable constraint.</td>
</tr>
<tr>
<td>conkey</td>
<td>int2[]</td>
<td>pg_attribute.attnum</td>
<td>If a table constraint (including foreign keys, but not constraint triggers), list of the constrained columns</td>
</tr>
<tr>
<td>confkey</td>
<td>int2[]</td>
<td>pg_attribute.attnum</td>
<td>If a foreign key, list of the referenced columns</td>
</tr>
<tr>
<td>conpfeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for PK = FK comparisons</td>
</tr>
<tr>
<td>conppeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for PK = PK comparisons</td>
</tr>
<tr>
<td>conffeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for FK = FK comparisons</td>
</tr>
<tr>
<td>conexclop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If an exclusion constraint, list of the per-column exclusion operators</td>
</tr>
<tr>
<td>conbin</td>
<td>pg_node_tree</td>
<td></td>
<td>If a check constraint, an internal representation of the expression</td>
</tr>
<tr>
<td>consrc</td>
<td>text</td>
<td></td>
<td>If a check constraint, a human-readable representation of the expression</td>
</tr>
</tbody>
</table>
In the case of an exclusion constraint, conkey is only useful for constraint elements that are simple column references. For other cases, a zero appears in conkey and the associated index must be consulted to discover the expression that is constrained. (conkey thus has the same contents as pg_index.indkey for the index.)

**Note**

consrc is not updated when referenced objects change; for example, it won't track renaming of columns. Rather than relying on this field, it's best to use pg_get_constraintdef() to extract the definition of a check constraint.

**Note**

pg_class.relchecks needs to agree with the number of check-constraint entries found in this table for each relation.

### 52.14. pg_conversion

The catalog pg_conversion describes encoding conversion functions. See CREATE CONVERSION for more information.

**Table 52.14. pg_conversion Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td>Row identifier</td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>conname</td>
<td>name</td>
<td>Conversion name</td>
<td>Conversion name (unique within a namespace)</td>
</tr>
<tr>
<td>connamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this conversion</td>
</tr>
<tr>
<td>conowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the conversion</td>
</tr>
<tr>
<td>conforencoding</td>
<td>int4</td>
<td></td>
<td>Source encoding ID</td>
</tr>
<tr>
<td>contoencoding</td>
<td>int4</td>
<td></td>
<td>Destination encoding ID</td>
</tr>
<tr>
<td>conproc</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Conversion function</td>
</tr>
<tr>
<td>condefault</td>
<td>bool</td>
<td></td>
<td>True if this is the default conversion</td>
</tr>
</tbody>
</table>

### 52.15. pg_database

The catalog pg_database stores information about the available databases. Databases are created with the CREATE DATABASE command. Consult Chapter 22 for details about the meaning of some of the parameters.

Unlike most system catalogs, pg_database is shared across all databases of a cluster: there is only one copy of pg_database per cluster, not one per database.
## Table 52.15. `pg_database` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td></td>
<td>Database name</td>
</tr>
<tr>
<td>datdba</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the database, usually the user who created it</td>
</tr>
<tr>
<td>encoding</td>
<td>int4</td>
<td></td>
<td>Character encoding for this database (pg_encoding_to_char() can translate this number to the encoding name)</td>
</tr>
<tr>
<td>datcollate</td>
<td>name</td>
<td></td>
<td>LC_COLLATE for this database</td>
</tr>
<tr>
<td>datctype</td>
<td>name</td>
<td></td>
<td>LC_CTYPE for this database</td>
</tr>
<tr>
<td>datistemplate</td>
<td>bool</td>
<td></td>
<td>If true, then this database can be cloned by any user with CREATEDB privileges; if false, then only superusers or the owner of the database can clone it.</td>
</tr>
<tr>
<td>datallowconn</td>
<td>bool</td>
<td></td>
<td>If false then no one can connect to this database. This is used to protect the template0 database from being altered.</td>
</tr>
<tr>
<td>datconnlimit</td>
<td>int4</td>
<td></td>
<td>Sets maximum number of concurrent connections that can be made to this database. -1 means no limit.</td>
</tr>
<tr>
<td>datlastsysoid</td>
<td>oid</td>
<td></td>
<td>Last system OID in the database; useful particularly to pg_dump</td>
</tr>
<tr>
<td>datfrozenxid</td>
<td>xid</td>
<td></td>
<td>All transaction IDs before this one have been replaced with a permanent (&quot;frozen&quot;) transaction ID in this database. This is used to track whether the database needs to be vacuumed in order to prevent transaction ID wraparound or to allow pg_xact to be shrunk. It is the minimum of the per-</td>
</tr>
</tbody>
</table>
### System Catalogs

#### Name | Type | References | Description
--- | --- | --- | ---
| datminmxid | xid | | All multixact IDs before this one have been replaced with a transaction ID in this database. This is used to track whether the database needs to be vacuumed in order to prevent multixact ID wraparound or to allow pg_multixact to be shrunk. It is the minimum of the per-table pg_class.relminmxid values.

| datatablespace | oid | pg_tablespace.oid | The default tablespace for the database. Within this database, all tables for which pg_class.reltablespace is zero will be stored in this tablespace; in particular, all the non-shared system catalogs will be there.

| datacl | aclitem[] | | Access privileges; see GRANT and REVOKE for details

---

## 52.16. pg_db_role_setting

The catalog pg_db_role_setting records the default values that have been set for run-time configuration variables, for each role and database combination.

Unlike most system catalogs, pg_db_role_setting is shared across all databases of a cluster: there is only one copy of pg_db_role_setting per cluster, not one per database.

### Table 52.16. pg_db_role_setting Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>setdatabase</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>The OID of the database the setting is applicable to, or zero if not database-specific</td>
</tr>
<tr>
<td>setrole</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>The OID of the role the setting is applicable to, or zero if not role-specific</td>
</tr>
<tr>
<td>setconfig</td>
<td>text[]</td>
<td></td>
<td>Defaults for run-time configuration variables</td>
</tr>
</tbody>
</table>
52.17. pg_default_acl

The catalog `pg_default_acl` stores initial privileges to be assigned to newly created objects.

Table 52.17. pg_default_acl Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>defaclrole</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>The OID of the role associated with this entry</td>
</tr>
<tr>
<td>defaclnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace associated with this entry, or 0 if none</td>
</tr>
<tr>
<td>defaclobjtype</td>
<td>char</td>
<td></td>
<td>Type of object this entry is for: r = relation (table, view), S = sequence, f = function, T = type, n = schema</td>
</tr>
<tr>
<td>defaclacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges that this type of object should have on creation</td>
</tr>
</tbody>
</table>

A `pg_default_acl` entry shows the initial privileges to be assigned to an object belonging to the indicated user. There are currently two types of entry: “global” entries with `defaclnamespace` = 0, and “per-schema” entries that reference a particular schema. If a global entry is present then it overrides the normal hard-wired default privileges for the object type. A per-schema entry, if present, represents privileges to be added to the global or hard-wired default privileges.

Note that when an ACL entry in another catalog is null, it is taken to represent the hard-wired default privileges for its object, not whatever might be in `pg_default_acl` at the moment. `pg_default_acl` is only consulted during object creation.

52.18. pg_depend

The catalog `pg_depend` records the dependency relationships between database objects. This information allows `DROP` commands to find which other objects must be dropped by `DROP CASCADE` or prevent dropping in the `DROP RESTRICT` case.

See also `pg_shdepend`, which performs a similar function for dependencies involving objects that are shared across a database cluster.

Table 52.18. pg_depend Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the dependent object is in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific dependent object</td>
</tr>
</tbody>
</table>
| objsubid | int4   |               | For a table column, this is the column number (the objid and classid refer to the table it-
In all cases, a `pg_depend` entry indicates that the referenced object cannot be dropped without also dropping the dependent object. However, there are several subflavors identified by `deptype`:

### DEPENDENCY_NORMAL (n)

A normal relationship between separately-created objects. The dependent object can be dropped without affecting the referenced object. The referenced object can only be dropped by specifying **CASCADE**, in which case the dependent object is dropped, too. Example: a table column has a normal dependency on its data type.

### DEPENDENCY_AUTO (a)

The dependent object can be dropped separately from the referenced object, and should be automatically dropped (regardless of **RESTRICT** or **CASCADE** mode) if the referenced object is dropped. Example: a named constraint on a table is made autodependent on the table, so that it will go away if the table is dropped.

### DEPENDENCY_INTERNAL (i)

The dependent object was created as part of creation of the referenced object, and is really just a part of its internal implementation. A **DROP** of the dependent object will be disallowed outright (we'll tell the user to issue a **DROP** against the referenced object, instead). A **DROP** of the referenced object will be propagated through to drop the dependent object whether **CASCADE** is specified or not. Example: a trigger that's created to enforce a foreign-key constraint is made internally dependent on the constraint's `pg_constraint` entry.

### DEPENDENCY_INTERNAL_AUTO (I)

The dependent object was created as part of creation of the referenced object, and is really just a part of its internal implementation. A **DROP** of the dependent object will be disallowed outright (we'll tell the user to issue a **DROP** against the referenced object, instead). While a regular internal dependency will prevent the dependent object from being dropped while any such dependencies remain, **DEPENDENCY_INTERNAL_AUTO** will allow such a drop as long as the object can be found by following any of such dependencies. Example: an index on a partition is made internal-auto-dependent on both the partition itself as well as on the index on the parent partitioned table; so the partition index is dropped together with either the partition it indexes, or with the parent index it is attached to.
DEPEN DENCY_EXTENSION (e)

The dependent object is a member of the extension that is the referenced object (see pg_extension). The dependent object can be dropped only via DROP EXTENSION on the referenced object. Functionally this dependency type acts the same as an internal dependency, but it's kept separate for clarity and to simplify pg_dump.

DEPEN DENCY_AUTO_EXTENSION (x)

The dependent object is not a member of the extension that is the referenced object (and so should not be ignored by pg_dump), but cannot function without it and should be dropped when the extension itself is. The dependent object may be dropped on its own as well.

DEPEN DENCY_PIN (p)

There is no dependent object; this type of entry is a signal that the system itself depends on the referenced object, and so that object must never be deleted. Entries of this type are created only by initdb. The columns for the dependent object contain zeroes.

Other dependency flavors might be needed in future.

52.19. pg_description

The catalog pg_description stores optional descriptions (comments) for each database object. Descriptions can be manipulated with the COMMENT command and viewed with psql's \d commands. Descriptions of many built-in system objects are provided in the initial contents of pg_description.

See also pg_shdescription, which performs a similar function for descriptions involving objects that are shared across a database cluster.

Table 52.19. pg_description Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this description pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a comment on a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>description</td>
<td>text</td>
<td></td>
<td>Arbitrary text that serves as the description of this object</td>
</tr>
</tbody>
</table>

52.20. pg_enum

The pg_enum catalog contains entries showing the values and labels for each enum type. The internal representation of a given enum value is actually the OID of its associated row in pg_enum.
### Table 52.20. pg_enum Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>enumtypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The OID of the pg_type entry owning this enum value</td>
</tr>
<tr>
<td>enumsortorder</td>
<td>float4</td>
<td></td>
<td>The sort position of this enum value within its enum type</td>
</tr>
<tr>
<td>enumlabel</td>
<td>name</td>
<td></td>
<td>The textual label for this enum value</td>
</tr>
</tbody>
</table>

The OIDs for pg_enum rows follow a special rule: even-numbered OIDs are guaranteed to be ordered in the same way as the sort ordering of their enum type. That is, if two even OIDs belong to the same enum type, the smaller OID must have the smaller enumsortorder value. Odd-numbered OID values need bear no relationship to the sort order. This rule allows the enum comparison routines to avoid catalog lookups in many common cases. The routines that create and alter enum types attempt to assign even OIDs to enum values whenever possible.

When an enum type is created, its members are assigned sort-order positions 1..n. But members added later might be given negative or fractional values of enumsortorder. The only requirement on these values is that they be correctly ordered and unique within each enum type.

#### 52.21. pg_event_trigger

The catalog pg_event_trigger stores event triggers. See Chapter 40 for more information.

### Table 52.21. pg_event_trigger Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>evtname</td>
<td>name</td>
<td></td>
<td>Trigger name (must be unique)</td>
</tr>
<tr>
<td>evtevent</td>
<td>name</td>
<td></td>
<td>Identifies the event for which this trigger fires</td>
</tr>
<tr>
<td>evtowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the event trigger</td>
</tr>
<tr>
<td>evtfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The function to be called</td>
</tr>
<tr>
<td>evtenabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the event trigger fires. O = trigger fires in “origin” and “local” modes, D = trigger is disabled, R = trigger fires in “replica” mode, A = trigger fires always.</td>
</tr>
<tr>
<td>evttags</td>
<td>text[]</td>
<td></td>
<td>Command tags for which this trigger will fire. If NULL, the firing of this trigger is not re-</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td><strong>stricted on the basis of</strong> the command tag.</td>
</tr>
</tbody>
</table>

## 52.22. pg_extension

The catalog `pg_extension` stores information about the installed extensions. See Section 38.16 for details about extensions.

### Table 52.22. pg_extension Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>extname</td>
<td>name</td>
<td></td>
<td>Name of the extension</td>
</tr>
<tr>
<td>extowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the extension</td>
</tr>
<tr>
<td>extnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Schema containing the extension's exported objects</td>
</tr>
<tr>
<td>extrelocatable</td>
<td>bool</td>
<td></td>
<td>True if extension can be relocated to another schema</td>
</tr>
<tr>
<td>extversion</td>
<td>text</td>
<td></td>
<td>Version name for the extension</td>
</tr>
<tr>
<td>extconfig</td>
<td>oid[]</td>
<td>pg_class.oid</td>
<td>Array of regclass OIDs for the extension's configuration table(s), or NULL if none</td>
</tr>
<tr>
<td>extcondition</td>
<td>text[]</td>
<td></td>
<td>Array of WHERE-clause filter conditions for the extension's configuration table(s), or NULL if none</td>
</tr>
</tbody>
</table>

Note that unlike most catalogs with a “namespace” column, `extnamespace` is not meant to imply that the extension belongs to that schema. Extension names are never schema-qualified. Rather, `extnamespace` indicates the schema that contains most or all of the extension's objects. If `extrelocatable` is true, then this schema must in fact contain all schema-qualifiable objects belonging to the extension.

## 52.23. pg_foreign_data_wrapper

The catalog `pg_foreign_data_wrapper` stores foreign-data wrapper definitions. A foreign-data wrapper is the mechanism by which external data, residing on foreign servers, is accessed.

### Table 52.23. pg_foreign_data_wrapper Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>fdwname</td>
<td>name</td>
<td></td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>fdwowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the foreign-data wrapper</td>
</tr>
<tr>
<td>fdwhandler</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>References a handler function that is responsible for supplying execution routines for the foreign-data wrapper. Zero if no handler is provided</td>
</tr>
<tr>
<td>fdwvalidator</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>References a validator function that is responsible for checking the validity of the options given to the foreign-data wrapper, as well as options for foreign servers and user mappings using the foreign-data wrapper. Zero if no validator is provided</td>
</tr>
<tr>
<td>fdwacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>fdwoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign-data wrapper specific options, as &quot;keyword=value&quot; strings</td>
</tr>
</tbody>
</table>

52.24. pg_foreign_server

The catalog pg_foreign_server stores foreign server definitions. A foreign server describes a source of external data, such as a remote server. Foreign servers are accessed via foreign-data wrappers.

Table 52.24. pg_foreign_server Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>srvname</td>
<td>name</td>
<td></td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>srvowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the foreign server</td>
</tr>
<tr>
<td>srvfdw</td>
<td>oid</td>
<td>pg_foreign_data_wrapper.oid</td>
<td>OID of the foreign-data wrapper of this foreign server</td>
</tr>
<tr>
<td>srvtype</td>
<td>text</td>
<td></td>
<td>Type of the server (optional)</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>srvversion</td>
<td>text</td>
<td></td>
<td>Version of the server (optional)</td>
</tr>
<tr>
<td>srvacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>srvoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign server specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

#### 52.25. pg_foreign_table

The catalog `pg_foreign_table` contains auxiliary information about foreign tables. A foreign table is primarily represented by a `pg_class` entry, just like a regular table. Its `pg_foreign_table` entry contains the information that is pertinent only to foreign tables and not any other kind of relation.

**Table 52.25. pg_foreign_table Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ftrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the <code>pg_class</code> entry for this foreign table</td>
</tr>
<tr>
<td>ftserver</td>
<td>oid</td>
<td>pg_foreign_server.oid</td>
<td>OID of the foreign server for this foreign table</td>
</tr>
<tr>
<td>ftoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign table options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

#### 52.26. pg_index

The catalog `pg_index` contains part of the information about indexes. The rest is mostly in `pg_class`.

**Table 52.26. pg_index Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>indexrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the <code>pg_class</code> entry for this index</td>
</tr>
<tr>
<td>indrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the <code>pg_class</code> entry for the table this index is for</td>
</tr>
<tr>
<td>indnatts</td>
<td>int2</td>
<td></td>
<td>The total number of columns in the index (duplicates <code>pg_class.relnatts</code>); this number includes both key and included attributes</td>
</tr>
<tr>
<td>indnkeyatts</td>
<td>int2</td>
<td></td>
<td>The number of key columns in the index, not counting any included columns, which are</td>
</tr>
</tbody>
</table>

2028
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>indisunique</td>
<td>bool</td>
<td></td>
<td>merely stored and do not participate in the index semantics</td>
</tr>
<tr>
<td>indisprimary</td>
<td>bool</td>
<td></td>
<td>If true, this is a unique index</td>
</tr>
<tr>
<td>indisexclusion</td>
<td>bool</td>
<td></td>
<td>If true, this index represents the primary key of the table (indisunique should always be true when this is true)</td>
</tr>
<tr>
<td>indimmediate</td>
<td>bool</td>
<td></td>
<td>If true, this index supports an exclusion constraint</td>
</tr>
<tr>
<td>indisclustered</td>
<td>bool</td>
<td></td>
<td>If true, the table was last clustered on this index</td>
</tr>
<tr>
<td>indisvalid</td>
<td>bool</td>
<td></td>
<td>If true, the index is currently valid for queries. False means the index is possibly incomplete: it must still be modified by INSERT/UPDATE operations, but it cannot safely be used for queries. If it is unique, the uniqueness property is not guaranteed true either.</td>
</tr>
<tr>
<td>indcheckxmin</td>
<td>bool</td>
<td></td>
<td>If true, queries must not use the index until the xmin of this pg_index row is below their TransactionXmin event horizon, because the table may contain broken HOT chains with incompatible rows that they can see</td>
</tr>
<tr>
<td>indisready</td>
<td>bool</td>
<td></td>
<td>If true, the index is currently ready for inserts. False means the index must be ignored by INSERT/UPDATE operations.</td>
</tr>
<tr>
<td>indislive</td>
<td>bool</td>
<td></td>
<td>If false, the index is in process of being dropped, and should be ignored for all purpos-</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>--------------------</td>
<td>------------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>indisreplident</td>
<td>bool</td>
<td></td>
<td>If true this index has been chosen as “repli-ca identity” using ALTER TABLE ... REPLICA IDENTITY USING INDEX ...</td>
</tr>
<tr>
<td>indkey</td>
<td>int2vector</td>
<td>pg_attribute.attnum</td>
<td>This is an array of indnatts values that indicate which table columns this index indexes. For example a value of 1 3 would mean that the first and the third table columns make up the index entries. Key columns come before non-key (included) columns. A zero in this array indicates that the corresponding index attribute is an expression over the table columns, rather than a simple column reference.</td>
</tr>
<tr>
<td>indcollation</td>
<td>oidvector</td>
<td>pg_collation.oid</td>
<td>For each column in the index key (indnkeyatts values), this contains the OID of the collation to use for the index, or zero if the column is not of a collatable data type.</td>
</tr>
<tr>
<td>indclass</td>
<td>oidvector</td>
<td>pg_opclass.oid</td>
<td>For each column in the index key (indnkeyatts values), this contains the OID of the operator class to use. See pg_opclass for details.</td>
</tr>
<tr>
<td>indoption</td>
<td>int2vector</td>
<td></td>
<td>This is an array of indnkeyatts values that store per-column flag bits. The meaning of the bits is defined by the index's access method.</td>
</tr>
<tr>
<td>indexprs</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression trees (in nodeToString() representation) for index attributes that are not simple column ref-</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>references. This is a list with one element for each zero entry in in-dkey. Null if all index attributes are simple references.</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>indpred</td>
<td>pg_node_tree</td>
<td>References</td>
<td>Expression tree (in nodeToString() representation) for partial index predicate. Null if not a partial index.</td>
</tr>
</tbody>
</table>

52.27. pg_inherits

The catalog pg_inherits records information about table inheritance hierarchies. There is one entry for each direct parent-child table relationship in the database. (Indirect inheritance can be determined by following chains of entries.)

Table 52.27. pg_inherits Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>inhrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the child table</td>
</tr>
<tr>
<td>inhparent</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the parent table</td>
</tr>
<tr>
<td>inhseqno</td>
<td>int4</td>
<td></td>
<td>If there is more than one direct parent for a child table (multiple inheritance), this number tells the order in which the inherited columns are to be arranged. The count starts at 1.</td>
</tr>
</tbody>
</table>

52.28. pg_init_privs

The catalog pg_init_privs records information about the initial privileges of objects in the system. There is one entry for each object in the database which has a non-default (non-NULL) initial set of privileges.

Objects can have initial privileges either by having those privileges set when the system is initialized (by initdb) or when the object is created during a CREATE EXTENSION and the extension script sets initial privileges using the GRANT system. Note that the system will automatically handle recording of the privileges during the extension script and that extension authors need only use the GRANT and REVOKE statements in their script to have the privileges recorded. The privtype column indicates if the initial privilege was set by initdb or during a CREATE EXTENSION command.

Objects which have initial privileges set by initdb will have entries where privtype is 'i', while objects which have initial privileges set by CREATE EXTENSION will have entries where privtype is 'e'.

2031
Table 52.28. pg_init_privs Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific object</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the object is in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>privtype</td>
<td>char</td>
<td></td>
<td>A code defining the type of initial privilege of this object; see text</td>
</tr>
<tr>
<td>initprivs</td>
<td>aclitem[]</td>
<td></td>
<td>The initial access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

52.29. pg_language

The catalog pg_language registers languages in which you can write functions or stored procedures. See CREATE LANGUAGE and Chapter 42 for more information about language handlers.

Table 52.29. pg_language Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>lanname</td>
<td>name</td>
<td></td>
<td>Name of the language</td>
</tr>
<tr>
<td>lanowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the language</td>
</tr>
<tr>
<td>lanispl</td>
<td>bool</td>
<td></td>
<td>This is false for internal languages (such as SQL) and true for user-defined languages. Currently, pg_dump still uses this to determine which languages need to be dumped, but this might be replaced by a different mechanism in the future.</td>
</tr>
<tr>
<td>lanpltrusted</td>
<td>bool</td>
<td></td>
<td>True if this is a trusted language, which means that it is believed not to grant access to anything outside the normal SQL execution environment. Only superusers can create functions in untrusted languages.</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>lanplcallfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>For noninternal languages this references the language handler, which is a special function that is responsible for executing all functions that are written in the particular language.</td>
</tr>
<tr>
<td>laninline</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>This references a function that is responsible for executing “inline” anonymous code blocks (DO blocks). Zero if inline blocks are not supported.</td>
</tr>
<tr>
<td>lanacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details.</td>
</tr>
</tbody>
</table>

**52.30. pg_largeobject**

The catalog `pg_largeobject` holds the data making up “large objects”. A large object is identified by an OID assigned when it is created. Each large object is broken into segments or “pages” small enough to be conveniently stored as rows in `pg_largeobject`. The amount of data per page is defined to be `LOBLKSIZE` (which is currently `BLCKSZ/4`, or typically 2 kB).

Prior to PostgreSQL 9.0, there was no permission structure associated with large objects. As a result, `pg_largeobject` was publicly readable and could be used to obtain the OIDs (and contents) of all large objects in the system. This is no longer the case; use `pg_largeobject_metadata` to obtain a list of large object OIDs.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>loid</td>
<td>oid</td>
<td>pg_largeobject_metadata.oid</td>
<td>Identifier of the large object that includes this page.</td>
</tr>
<tr>
<td>pageno</td>
<td>int4</td>
<td></td>
<td>Page number of this page within its large object (counting from zero)</td>
</tr>
<tr>
<td>data</td>
<td>bytea</td>
<td></td>
<td>Actual data stored in the large object. This will never be more than <code>LOBLKSIZE</code> bytes and might be less.</td>
</tr>
</tbody>
</table>
Each row of `pg_largeobject` holds data for one page of a large object, beginning at byte offset `pageno * LOBLKSIZE` within the object. The implementation allows sparse storage: pages might be missing, and might be shorter than `LOBLKSIZE` bytes even if they are not the last page of the object. Missing regions within a large object read as zeroes.

### 52.31. `pg_largeobject_metadata`

The catalog `pg_largeobject_metadata` holds metadata associated with large objects. The actual large object data is stored in `pg_largeobject`.

**Table 52.31. `pg_largeobject_metadata` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>lomowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the large object</td>
</tr>
<tr>
<td>lomacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

### 52.32. `pg_namespace`

The catalog `pg_namespace` stores namespaces. A namespace is the structure underlying SQL schemas: each namespace can have a separate collection of relations, types, etc. without name conflicts.

**Table 52.32. `pg_namespace` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>nspname</td>
<td>name</td>
<td></td>
<td>Name of the namespace</td>
</tr>
<tr>
<td>nspowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the namespace</td>
</tr>
<tr>
<td>nspacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

### 52.33. `pg_opclass`

The catalog `pg_opclass` defines index access method operator classes. Each operator class defines semantics for index columns of a particular data type and a particular index access method. An operator class essentially specifies that a particular operator family is applicable to a particular indexable column data type. The set of operators from the family that are actually usable with the indexed column are whichever ones accept the column’s data type as their left-hand input.

Operator classes are described at length in Section 38.15.
Table 52.33. pg_opclass Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>opcmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator class is for</td>
</tr>
<tr>
<td>opcname</td>
<td>name</td>
<td></td>
<td>Name of this operator class</td>
</tr>
<tr>
<td>opcnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Namespace of this operator class</td>
</tr>
<tr>
<td>opcowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator class</td>
</tr>
<tr>
<td>opcfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>Operator family containing the operator class</td>
</tr>
<tr>
<td>opcintype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type that the operator class indexes</td>
</tr>
<tr>
<td>opcdefault</td>
<td>bool</td>
<td></td>
<td>True if this operator class is the default for opcintype</td>
</tr>
<tr>
<td>opckeytype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of data stored in index, or zero if same as opcintype</td>
</tr>
</tbody>
</table>

An operator class’s opcmethod must match the opfmethod of its containing operator family. Also, there must be no more than one pg_opclass row having opcdefault true for any given combination of opcmethod and opcintype.

52.34. pg_operator

The catalog pg_operator stores information about operators. See CREATE OPERATOR and Section 38.13 for more information.

Table 52.34. pg_operator Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>oprname</td>
<td>name</td>
<td></td>
<td>Name of the operator</td>
</tr>
<tr>
<td>oprnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this operator</td>
</tr>
<tr>
<td>oprowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator</td>
</tr>
<tr>
<td>oprkind</td>
<td>char</td>
<td></td>
<td>b = infix (&quot;both&quot;), l = prefix (&quot;left&quot;), r = postfix (&quot;right&quot;)</td>
</tr>
<tr>
<td>oprcanmerge</td>
<td>bool</td>
<td></td>
<td>This operator supports merge joins</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oprcanhash</td>
<td>bool</td>
<td></td>
<td>This operator supports hash joins</td>
</tr>
<tr>
<td>oprleft</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the left operand</td>
</tr>
<tr>
<td>oprright</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the right operand</td>
</tr>
<tr>
<td>oprresult</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the result</td>
</tr>
<tr>
<td>oprcom</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Commutator of this operator, if any</td>
</tr>
<tr>
<td>oprnegate</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Negator of this operator, if any</td>
</tr>
<tr>
<td>oprcode</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Function that implements this operator</td>
</tr>
<tr>
<td>oprrest</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Restriction selectivity estimation function for this operator</td>
</tr>
<tr>
<td>oprjoin</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Join selectivity estimation function for this operator</td>
</tr>
</tbody>
</table>

Unused column contain zeroes. For example, oprleft is zero for a prefix operator.

52.35. pg_opfamily

The catalog pg_opfamily defines operator families. Each operator family is a collection of operators and associated support routines that implement the semantics specified for a particular index access method. Furthermore, the operators in a family are all “compatible”, in a way that is specified by the access method. The operator family concept allows cross-data-type operators to be used with indexes and to be reasoned about using knowledge of access method semantics.

Operator families are described at length in Section 38.15.

Table 52.35. pg_opfamily Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>opfmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator family is for</td>
</tr>
<tr>
<td>opfname</td>
<td>name</td>
<td></td>
<td>Name of this operator family</td>
</tr>
<tr>
<td>opfnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Namespace of this operator family</td>
</tr>
<tr>
<td>opfowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator family</td>
</tr>
</tbody>
</table>

The majority of the information defining an operator family is not in its pg_opfamily row, but in the associated rows in pg_amop, pg_amproc, and pg_opclass.

52.36. pg_partitioned_table
The catalog `pg_partitioned_table` stores information about how tables are partitioned.

### Table 52.36. `pg_partitioned_table` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>partrelid</code></td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The OID of the <code>pg_class</code> entry for this partitioned table</td>
</tr>
<tr>
<td><code>partstrat</code></td>
<td>char</td>
<td></td>
<td>Partitioning strategy; h = hash partitioned table, l = list partitioned table, r = range partitioned table</td>
</tr>
<tr>
<td><code>partnatts</code></td>
<td>int2</td>
<td></td>
<td>The number of columns in partition key</td>
</tr>
<tr>
<td><code>partdefid</code></td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The OID of the <code>pg_class</code> entry for the default partition of this partitioned table, or zero if this partitioned table does not have a default partition.</td>
</tr>
<tr>
<td><code>partattrs</code></td>
<td>int2vector</td>
<td><code>pg_attribute.attnum</code></td>
<td>This is an array of <code>partnatts</code> values that indicate which table columns are part of the partition key. For example, a value of 1, 3 would mean that the first and the third table columns make up the partition key. A zero in this array indicates that the corresponding partition key column is an expression, rather than a simple column reference.</td>
</tr>
<tr>
<td><code>partclass</code></td>
<td>oidvector</td>
<td><code>pg_opclass.oid</code></td>
<td>For each column in the partition key, this contains the OID of the operator class to use. See <code>pg_opclass</code> for details.</td>
</tr>
<tr>
<td><code>partcollation</code></td>
<td>oidvector</td>
<td><code>pg_opclass.oid</code></td>
<td>For each column in the partition key, this contains the OID of the collation to use for partitioning, or zero if the column is not of a collatable data type.</td>
</tr>
<tr>
<td><code>partexprs</code></td>
<td><code>pg_node_tree</code></td>
<td></td>
<td>Expression trees (in <code>nodeToString()</code> representation) for partition key columns that</td>
</tr>
</tbody>
</table>
52.37. pg_pltemplate

The catalog pg_pltemplate stores “template” information for procedural languages. A template for a language allows the language to be created in a particular database by a simple CREATE LANGUAGE command, with no need to specify implementation details.

Unlike most system catalogs, pg_pltemplate is shared across all databases of a cluster: there is only one copy of pg_pltemplate per cluster, not one per database. This allows the information to be accessible in each database as it is needed.

Table 52.37. pg_pltemplate Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>tmplname</td>
<td>name</td>
<td>Name of the language this template is for</td>
</tr>
<tr>
<td>tmppltrusted</td>
<td>boolean</td>
<td>True if language is considered trusted</td>
</tr>
<tr>
<td>tmpldbcreate</td>
<td>boolean</td>
<td>True if language may be created by a database owner</td>
</tr>
<tr>
<td>tmpplhandler</td>
<td>text</td>
<td>Name of call handler function</td>
</tr>
<tr>
<td>tmplinline</td>
<td>text</td>
<td>Name of anonymous-block handler function, or null if none</td>
</tr>
<tr>
<td>tmpplvalidator</td>
<td>text</td>
<td>Name of validator function, or null if none</td>
</tr>
<tr>
<td>tmpllibrary</td>
<td>text</td>
<td>Path of shared library that implements language</td>
</tr>
<tr>
<td>tmplacl</td>
<td>aclitem[]</td>
<td>Access privileges for template (not actually used)</td>
</tr>
</tbody>
</table>

There are not currently any commands that manipulate procedural language templates; to change the built-in information, a superuser must modify the table using ordinary INSERT, DELETE, or UPDATE commands.

Note

It is likely that pg_pltemplate will be removed in some future release of PostgreSQL, in favor of keeping this knowledge about procedural languages in their respective extension installation scripts.

52.38. pg_policy

The catalog pg_policy stores row level security policies for tables. A policy includes the kind of command that it applies to (possibly all commands), the roles that it applies to, the expression to be
added as a security-barrier qualification to queries that include the table, and the expression to be added as a WITH CHECK option for queries that attempt to add new records to the table.

<table>
<thead>
<tr>
<th>Table 52.38. pg_policy Columns</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Name</strong></td>
</tr>
<tr>
<td>polname</td>
</tr>
<tr>
<td>polrelid</td>
</tr>
<tr>
<td>polcmd</td>
</tr>
<tr>
<td>polpermissive</td>
</tr>
<tr>
<td>polroles</td>
</tr>
<tr>
<td>polqual</td>
</tr>
<tr>
<td>polwithcheck</td>
</tr>
</tbody>
</table>

**Note**

Policies stored in pg_policy are applied only when pg_class.relrowsecurity is set for their table.

52.39. pg_proc

The catalog pg_proc stores information about functions, procedures, aggregate functions, and window functions (collectively also known as routines). See CREATE FUNCTION, CREATE PROCEDURE, and Section 38.3 for more information.

If prokind indicates that the entry is for an aggregate function, there should be a matching row in pg_aggregate.

<table>
<thead>
<tr>
<th>Table 52.39. pg_proc Columns</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Name</strong></td>
</tr>
<tr>
<td>oid</td>
</tr>
<tr>
<td>proname</td>
</tr>
<tr>
<td>pronamespace</td>
</tr>
<tr>
<td>Name</td>
</tr>
<tr>
<td>-----------------</td>
</tr>
<tr>
<td>proowner</td>
</tr>
<tr>
<td>prolang</td>
</tr>
<tr>
<td>procost</td>
</tr>
<tr>
<td>prorows</td>
</tr>
<tr>
<td>provariadic</td>
</tr>
<tr>
<td>protransform</td>
</tr>
<tr>
<td>prokind</td>
</tr>
<tr>
<td>prosecdef</td>
</tr>
<tr>
<td>proleakproof</td>
</tr>
<tr>
<td>proisstrict</td>
</tr>
<tr>
<td>proretset</td>
</tr>
<tr>
<td>provolatile</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th><strong>Name</strong></th>
<th><strong>Type</strong></th>
<th><strong>References</strong></th>
<th><strong>Description</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>its input arguments, or is affected by outside factors. It is i for “immutable” functions, which always deliver the same result for the same inputs. It is s for “stable” functions, whose results (for fixed inputs) do not change within a scan. It is v for “volatile” functions, whose results might change at any time. (Use v also for functions with side-effects, so that calls to them cannot get optimized away.)</td>
</tr>
<tr>
<td>proparallel</td>
<td>char</td>
<td></td>
<td>proparallel tells whether the function can be safely run in parallel mode. It is s for functions which are safe to run in parallel mode without restriction. It is r for functions which can be run in parallel mode, but their execution is restricted to the parallel group leader; parallel worker processes cannot invoke these functions. It is u for functions which are unsafe in parallel mode; the presence of such a function forces a serial execution plan.</td>
</tr>
<tr>
<td>pronargs</td>
<td>int2</td>
<td></td>
<td>Number of input arguments</td>
</tr>
<tr>
<td>pronargdefaults</td>
<td>int2</td>
<td></td>
<td>Number of arguments that have defaults</td>
</tr>
<tr>
<td>prorettype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the return value</td>
</tr>
<tr>
<td>proargtypes</td>
<td>oidvector</td>
<td>pg_type.oid</td>
<td>An array with the data types of the function arguments. This includes only input arguments (including INOUT and VARIADIC arguments), and thus represents the call signature of the function.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>---------------</td>
<td>------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>proallargtypes</td>
<td>oid[]</td>
<td>pg_type.oid</td>
<td>An array with the data types of the function arguments. This includes all arguments (including OUT and INOUT arguments); however, if all the arguments are IN arguments, this field will be null. Note that subscripting is 1-based, whereas for historical reasons proargtypes is subscripted from 0.</td>
</tr>
<tr>
<td>proargmodes</td>
<td>char[]</td>
<td></td>
<td>An array with the modes of the function arguments, encoded as i for IN arguments, o for OUT arguments, b for INOUT arguments, v for VARIADIC arguments, t for TABLE arguments. If all the arguments are IN arguments, this field will be null. Note that subscripts correspond to positions of proallargtypes not proargtypes.</td>
</tr>
<tr>
<td>proargnames</td>
<td>text[]</td>
<td></td>
<td>An array with the names of the function arguments. Arguments without a name are set to empty strings in the array. If none of the arguments have a name, this field will be null. Note that subscripts correspond to positions of proallargtypes not proargtypes.</td>
</tr>
<tr>
<td>proargdefaults</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression trees (in nodeToString() representation) for default values. This is a list with proargdefaults elements corresponding to the last N input arguments (i.e., the last Nproargtypes positions). If none of the arguments have defaults, this field will be null.</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>protrftypes</td>
<td>oid[]</td>
<td></td>
<td>Data type OIDs for which to apply transforms.</td>
</tr>
<tr>
<td>prosrc</td>
<td>text</td>
<td></td>
<td>This tells the function handler how to invoke the function. It might be</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>the actual source code of the function for interpreted languages, a link</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>symbol, a file name, or just about anything else, depending on the</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>implementation language/call convention.</td>
</tr>
<tr>
<td>probin</td>
<td>text</td>
<td></td>
<td>Additional information about how to invoke the function. Again, the</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>interpretation is language-specific.</td>
</tr>
<tr>
<td>proconfig</td>
<td>text[]</td>
<td></td>
<td>Function’s local settings for run-time configuration variables</td>
</tr>
<tr>
<td>proacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

For compiled functions, both built-in and dynamically loaded, prosrc contains the function’s C-language name (link symbol). For all other currently-known language types, prosrc contains the function’s source text. probin is unused except for dynamically-loaded C functions, for which it gives the name of the shared library file containing the function.

### 52.40. pg_publication

The catalog pg_publication contains all publications created in the database. For more on publications see Section 31.1.

#### Table 52.40. pg_publication Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>pubname</td>
<td>name</td>
<td></td>
<td>Name of the publication</td>
</tr>
<tr>
<td>pubowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the publication</td>
</tr>
<tr>
<td>puballtables</td>
<td>bool</td>
<td></td>
<td>If true, this publication automatically includes all tables in the database,</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>including any that will be created in the future.</td>
</tr>
</tbody>
</table>
### 52.41. `pg_publication_rel`

The catalog `pg_publication_rel` contains the mapping between relations and publications in the database. This is a many-to-many mapping. See also Section 52.78 for a more user-friendly view of this information.

**Table 52.41. pg_publication_rel Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>prpubid</td>
<td>oid</td>
<td><code>pg_publication.oid</code></td>
<td>Reference to publication</td>
</tr>
<tr>
<td>prrelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>Reference to relation</td>
</tr>
</tbody>
</table>

### 52.42. `pg_range`

The catalog `pg_range` stores information about range types. This is in addition to the types’ entries in `pg_type`.

**Table 52.42. pg_range Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rntypid</td>
<td>oid</td>
<td><code>pg_type.oid</code></td>
<td>OID of the range type</td>
</tr>
<tr>
<td>rnsubtype</td>
<td>oid</td>
<td><code>pg_type.oid</code></td>
<td>OID of the element type (subtype) of this range type</td>
</tr>
<tr>
<td>rngcollation</td>
<td>oid</td>
<td><code>pg_collation.oid</code></td>
<td>OID of the collation used for range comparisons, or 0 if none</td>
</tr>
<tr>
<td>rnsubopc</td>
<td>oid</td>
<td><code>pg_opclass.oid</code></td>
<td>OID of the subtype's operator class used for range comparisons</td>
</tr>
<tr>
<td>rngcanonical</td>
<td>regproc</td>
<td><code>pg_proc.oid</code></td>
<td>OID of the function to convert a range value into canonical form, or 0 if none</td>
</tr>
<tr>
<td>rnsubdiff</td>
<td>regproc</td>
<td><code>pg_proc.oid</code></td>
<td>OID of the function to return the difference be-</td>
</tr>
</tbody>
</table>
rngsubopc (plus rngcollation, if the element type is collatable) determines the sort ordering used by the range type. rngcanonical is used when the element type is discrete. rngsubdiff is optional but should be supplied to improve performance of GiST indexes on the range type.

### 52.43. pg_replication_origin

The pg_replication_origin catalog contains all replication origins created. For more on replication origins see Chapter 50.

Unlike most system catalogs, pg_replication_origin is shared across all databases of a cluster: there is only one copy of pg_replication_origin per cluster, not one per database.

**Table 52.43. pg_replication_origin Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>roident</td>
<td>Oid</td>
<td></td>
<td>A unique, cluster-wide identifier for the replication origin. Should never leave the system.</td>
</tr>
<tr>
<td>roname</td>
<td>text</td>
<td></td>
<td>The external, user defined, name of a replication origin.</td>
</tr>
</tbody>
</table>

### 52.44. pg_rewrite

The catalog pg_rewrite stores rewrite rules for tables and views.

**Table 52.44. pg_rewrite Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>rulename</td>
<td>name</td>
<td></td>
<td>Rule name</td>
</tr>
<tr>
<td>ev_class</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this rule is for</td>
</tr>
<tr>
<td>ev_type</td>
<td>char</td>
<td></td>
<td>Event type that the rule is for: 1 = SELECT, 2 = UPDATE, 3 = INSERT, 4 = DELETE</td>
</tr>
<tr>
<td>ev_enabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the rule fires. O = rule fires in “origin” and “local” modes, D = rule is disabled, R = rule fires in “replica” mode, A = rule fires always.</td>
</tr>
<tr>
<td>is_instead</td>
<td>bool</td>
<td></td>
<td>True if the rule is an INSTEAD rule</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ev_qual</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression tree (in the form of a nodeToString() representation) for the rule's qualifying condition</td>
</tr>
<tr>
<td>ev_action</td>
<td>pg_node_tree</td>
<td></td>
<td>Query tree (in the form of a nodeToString() representation) for the rule's action</td>
</tr>
</tbody>
</table>

Note

`pg_class.relhasrules` must be true if a table has any rules in this catalog.

52.45. pg_seclabel

The catalog `pg_seclabel` stores security labels on database objects. Security labels can be manipulated with the SECURITY LABEL command. For an easier way to view security labels, see Section 52.83.

See also `pg_shseclabel`, which performs a similar function for security labels of database objects that are shared across a database cluster.

Table 52.45. pg_seclabel Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a security label on a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td></td>
<td>The label provider associated with this label.</td>
</tr>
<tr>
<td>label</td>
<td>text</td>
<td></td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

52.46. pg_sequence

The catalog `pg_sequence` contains information about sequences. Some of the information about sequences, such as the name and the schema, is in `pg_class`. 

2046
### Table 52.46. pg_sequence Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>seqrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the pg_class entry for this sequence</td>
</tr>
<tr>
<td>seqtypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the sequence</td>
</tr>
<tr>
<td>seqstart</td>
<td>int8</td>
<td></td>
<td>Start value of the sequence</td>
</tr>
<tr>
<td>seqincrement</td>
<td>int8</td>
<td></td>
<td>Increment value of the sequence</td>
</tr>
<tr>
<td>seqmax</td>
<td>int8</td>
<td></td>
<td>Maximum value of the sequence</td>
</tr>
<tr>
<td>seqmin</td>
<td>int8</td>
<td></td>
<td>Minimum value of the sequence</td>
</tr>
<tr>
<td>seqcache</td>
<td>int8</td>
<td></td>
<td>Cache size of the sequence</td>
</tr>
<tr>
<td>seqcycle</td>
<td>bool</td>
<td></td>
<td>Whether the sequence cycles</td>
</tr>
</tbody>
</table>

### 52.47. pg_shdepend

The catalog pg_shdepend records the dependency relationships between database objects and shared objects, such as roles. This information allows PostgreSQL to ensure that those objects are unreferenced before attempting to delete them.

See also pg_depend, which performs a similar function for dependencies involving objects within a single database.

Unlike most system catalogs, pg_shdepend is shared across all databases of a cluster: there is only one copy of pg_shdepend per cluster, not one per database.

### Table 52.47. pg_shdepend Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>dbid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>The OID of the database the dependent object is in, or zero for a shared object</td>
</tr>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the dependent object is in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific dependent object</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the objid and classid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>refclassid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the referenced object is in (must be a shared catalog)</td>
</tr>
<tr>
<td>refobjid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific referenced object</td>
</tr>
<tr>
<td>deptype</td>
<td>char</td>
<td></td>
<td>A code defining the specific semantics of this dependency relationship; see text</td>
</tr>
</tbody>
</table>

In all cases, a pg_shdepend entry indicates that the referenced object cannot be dropped without also dropping the dependent object. However, there are several subflavors identified by deptype:

**SHARED_DEPENDENCY_OWNER (o)**

The referenced object (which must be a role) is the owner of the dependent object.

**SHARED_DEPENDENCY_ACL (a)**

The referenced object (which must be a role) is mentioned in the ACL (access control list, i.e., privileges list) of the dependent object. (A SHARED_DEPENDENCY_ACL entry is not made for the owner of the object, since the owner will have a SHARED_DEPENDENCY_OWNER entry anyway.)

**SHARED_DEPENDENCY_POLICY (r)**

The referenced object (which must be a role) is mentioned as the target of a dependent policy object.

**SHARED_DEPENDENCY_PIN (p)**

There is no dependent object; this type of entry is a signal that the system itself depends on the referenced object, and so that object must never be deleted. Entries of this type are created only by initdb. The columns for the dependent object contain zeroes.

Other dependency flavors might be needed in future. Note in particular that the current definition only supports roles as referenced objects.

### 52.48. pg_shdescription

The catalog pg_shdescription stores optional descriptions (comments) for shared database objects. Descriptions can be manipulated with the COMMENT command and viewed with psql's \d commands.

See also pg_description, which performs a similar function for descriptions involving objects within a single database.

Unlike most system catalogs, pg_shdescription is shared across all databases of a cluster: there is only one copy of pg_shdescription per cluster, not one per database.

#### Table 52.48. pg_shdescription Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this description pertains to</td>
</tr>
</tbody>
</table>
### 52.49. `pg_shseclabel`

The catalog `pg_shseclabel` stores security labels on shared database objects. Security labels can be manipulated with the `SECURITY LABEL` command. For an easier way to view security labels, see Section 52.83.

See also `pg_seclabel`, which performs a similar function for security labels involving objects within a single database.

Unlike most system catalogs, `pg_shseclabel` is shared across all databases of a cluster: there is only one copy of `pg_shseclabel` per cluster, not one per database.

### Table 52.49. `pg_shseclabel` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td></td>
<td>The label provider associated with this label.</td>
</tr>
<tr>
<td>label</td>
<td>text</td>
<td></td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

### 52.50. `pg_statistic`

The catalog `pg_statistic` stores statistical data about the contents of the database. Entries are created by `ANALYZE` and subsequently used by the query planner. Note that all the statistical data is inherently approximate, even assuming that it is up-to-date.

Normally there is one entry, with `stainherit = false`, for each table column that has been analyzed. If the table has inheritance children, a second entry with `stainherit = true` is also created. This row represents the column's statistics over the inheritance tree, i.e., statistics for the data you'd see with `SELECT column FROM table*`, whereas the `stainherit = false` row represents the results of `SELECT column FROM ONLY table`.

`pg_statistic` also stores statistical data about the values of index expressions. These are described as if they were actual data columns; in particular, `starelid` references the index. No entry is made for an ordinary non-expression index column, however, since it would be redundant with the entry for the underlying table column. Currently, entries for index expressions always have `stainherit = false`.

Since different kinds of statistics might be appropriate for different kinds of data, `pg_statistic` is designed not to assume very much about what sort of statistics it stores. Only extremely general
statistics (such as nullness) are given dedicated columns in `pg_statistic`. Everything else is stored in “slots”, which are groups of associated columns whose content is identified by a code number in one of the slot’s columns. For more information see `src/include/catalog/pg_statistic.h`.

`pg_statistic` should not be readable by the public, since even statistical information about a table’s contents might be considered sensitive. (Example: minimum and maximum values of a salary column might be quite interesting.) `pg_stats` is a publicly readable view on `pg_statistic` that only exposes information about those tables that are readable by the current user.

### Table 52.50. `pg_statistic` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>starelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The table or index that the described column belongs to</td>
</tr>
<tr>
<td>staattnum</td>
<td>int2</td>
<td><code>pg_attribute.attnum</code></td>
<td>The number of the described column</td>
</tr>
<tr>
<td>stainherit</td>
<td>bool</td>
<td></td>
<td>If true, the stats include inheritance child columns, not just the values in the specified relation</td>
</tr>
<tr>
<td>stanullfrac</td>
<td>float4</td>
<td></td>
<td>The fraction of the column’s entries that are null</td>
</tr>
<tr>
<td>stawidth</td>
<td>int4</td>
<td></td>
<td>The average stored width, in bytes, of non-null entries</td>
</tr>
<tr>
<td>stadistinct</td>
<td>float4</td>
<td></td>
<td>The number of distinct nonnull data values in the column. A value greater than zero is the actual number of distinct values. A value less than zero is the negative of a multiplier for the number of rows in the table; for example, a column in which about 80% of the values are nonnull and each nonnull value appears about twice on average could be represented by <code>stadistinct = -0.4</code>. A zero value means the number of distinct values is unknown.</td>
</tr>
<tr>
<td>stakindN</td>
<td>int2</td>
<td></td>
<td>A code number indicating the kind of statistics stored in the Nth “slot” of the <code>pg_statistic</code> row.</td>
</tr>
<tr>
<td>staopN</td>
<td>oid</td>
<td><code>pg_operator.oid</code></td>
<td>An operator used to derive the statistics stored</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stanumbersN</td>
<td>float4[]</td>
<td>null</td>
<td>Numerical statistics of the appropriate kind for the Nth “slot”, or null if the slot kind does not involve numerical values</td>
</tr>
<tr>
<td>stavaluesN</td>
<td>anyarray</td>
<td>null</td>
<td>Column data values of the appropriate kind for the Nth “slot”, or null if the slot kind does not store any data values. Each array’s element values are actually of the specific column’s data type, or a related type such as an array’s element type, so there is no way to define these columns’ type more specifically than anyarray.</td>
</tr>
</tbody>
</table>

52.51. pg_statistic_ext

The catalog pg_statistic_ext holds extended planner statistics. Each row in this catalog corresponds to a statistics object created with CREATE STATISTICS.

Table 52.51. pg_statistic_ext Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stxrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>Table containing the columns described by this object</td>
</tr>
<tr>
<td>stxname</td>
<td>name</td>
<td></td>
<td>Name of the statistics object</td>
</tr>
<tr>
<td>stxnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this statistics object</td>
</tr>
<tr>
<td>stxowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the statistics object</td>
</tr>
<tr>
<td>stxkeys</td>
<td>int2vector</td>
<td>pg_attribute.attnum</td>
<td>An array of attribute numbers, indicating which table columns are covered by this statistics object; for example a value of 1 3 would mean that the first and the third table columns are covered</td>
</tr>
</tbody>
</table>
System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stxkind</td>
<td>char[]</td>
<td></td>
<td>An array containing codes for the enabled statistic kinds; valid values are: d for n-distinct statistics, f for functional dependency statistics</td>
</tr>
<tr>
<td>stxndistinct</td>
<td>pg_ndistinct</td>
<td></td>
<td>N-distinct counts, serialized as pg_ndistinct type</td>
</tr>
<tr>
<td>stxdependencies</td>
<td>pg_dependencies</td>
<td></td>
<td>Functional dependency statistics, serialized as pg_dependencies type</td>
</tr>
</tbody>
</table>

The stxkind field is filled at creation of the statistics object, indicating which statistic type(s) are desired. The fields after it are initially NULL and are filled only when the corresponding statistic has been computed by ANALYZE.

52.52. pg_subscription

The catalog pg_subscription contains all existing logical replication subscriptions. For more information about logical replication see Chapter 31.

Unlike most system catalogs, pg_subscription is shared across all databases of a cluster: there is only one copy of pg_subscription per cluster, not one per database.

Access to the column subconninfo is revoked from normal users, because it could contain plain-text passwords.

Table 52.52. pg_subscription Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>subdbid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>OID of the database which the subscription resides in</td>
</tr>
<tr>
<td>subname</td>
<td>name</td>
<td></td>
<td>Name of the subscription</td>
</tr>
<tr>
<td>subowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the subscription</td>
</tr>
<tr>
<td>subenabled</td>
<td>bool</td>
<td></td>
<td>If true, the subscription is enabled and should be replicating.</td>
</tr>
<tr>
<td>subsynccommit</td>
<td>text</td>
<td></td>
<td>Contains the value of the synchronous_commit setting for the subscription workers.</td>
</tr>
<tr>
<td>subconninfo</td>
<td>text</td>
<td></td>
<td>Connection string to the upstream database</td>
</tr>
</tbody>
</table>
### System Catalogs

#### 52.53. pg_subscription_rel

The catalog `pg_subscription_rel` contains the state for each replicated relation in each subscription. This is a many-to-many mapping.

This catalog only contains tables known to the subscription after running either `CREATE SUBSCRIPTION` or `ALTER SUBSCRIPTION ... REFRESH PUBLICATION`.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>srssubid</td>
<td>oid</td>
<td><code>pg_subscription.oid</code></td>
<td>Reference to subscription</td>
</tr>
<tr>
<td>srrelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>Reference to relation</td>
</tr>
<tr>
<td>srsstate</td>
<td>char</td>
<td></td>
<td>State code: i = initialize, d = data is being copied, s = synchronized, r = ready (normal replication)</td>
</tr>
<tr>
<td>srsblsn</td>
<td><code>pg_lsn</code></td>
<td></td>
<td>End LSN for s and r states.</td>
</tr>
</tbody>
</table>

#### 52.54. pg_tablespace

The catalog `pg_tablespace` stores information about the available tablespaces. Tables can be placed in particular tablespaces to aid administration of disk layout.

Unlike most system catalogs, `pg_tablespace` is shared across all databases of a cluster: there is only one copy of `pg_tablespace` per cluster, not one per database.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>spcname</td>
<td>name</td>
<td></td>
<td>Tablespace name</td>
</tr>
<tr>
<td>spcowner</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>Owner of the tablespace, usually the user who created it</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>------------</td>
<td>------------</td>
<td>------------------------------------------</td>
</tr>
<tr>
<td>spcacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>spcoptions</td>
<td>text[]</td>
<td></td>
<td>Tablespace-level options, as &quot;keyword=value&quot; strings</td>
</tr>
</tbody>
</table>

### 52.55. pg_transform

The catalog pg_transform stores information about transforms, which are a mechanism to adapt data types to procedural languages. See CREATE TRANSFORM for more information.

**Table 52.55. pg_transform Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trftype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the data type this transform is for</td>
</tr>
<tr>
<td>trflang</td>
<td>oid</td>
<td>pg_language.oid</td>
<td>OID of the language this transform is for</td>
</tr>
<tr>
<td>trffromsql</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use when converting the data type for input to the procedural language (e.g., function parameters). Zero is stored if this operation is not supported.</td>
</tr>
<tr>
<td>trftosql</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use when converting output from the procedural language (e.g., return values) to the data type. Zero is stored if this operation is not supported.</td>
</tr>
</tbody>
</table>

### 52.56. pg_trigger

The catalog pg_trigger stores triggers on tables and views. See CREATE TRIGGER for more information.

**Table 52.56. pg_trigger Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>tgrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this trigger is on</td>
</tr>
<tr>
<td>tname</td>
<td>name</td>
<td></td>
<td>Trigger name (must be unique among triggers of same table)</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>---------------</td>
<td>---------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>tgfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The function to be called</td>
</tr>
<tr>
<td>tgtype</td>
<td>int2</td>
<td></td>
<td>Bit mask identifying trigger firing conditions</td>
</tr>
<tr>
<td>tgenabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the trigger fires.</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>O = trigger fires in “origin” and “local” modes,</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>D = trigger is disabled,</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>R = trigger fires in “repli-</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>cala” mode, A = trigger fires always.</td>
</tr>
<tr>
<td>tgisinternal</td>
<td>bool</td>
<td></td>
<td>True if trigger is internally generated (usually, to enforce the</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>constraint identified by tgconstraint)</td>
</tr>
<tr>
<td>tgconstrrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table referenced by a referential integrity constraint</td>
</tr>
<tr>
<td>tgconstrindid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The index supporting a unique, primary key, referential integrity, or</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>exclusion constraint</td>
</tr>
<tr>
<td>tgconstraint</td>
<td>oid</td>
<td>pg_constraint.oid</td>
<td>The pg_constraint entry associated with the trigger, if any</td>
</tr>
<tr>
<td>tgdeferrable</td>
<td>bool</td>
<td></td>
<td>True if constraint trigger is deferrable</td>
</tr>
<tr>
<td>tginitdeferred</td>
<td>bool</td>
<td></td>
<td>True if constraint trigger is initially deferred</td>
</tr>
<tr>
<td>tgnargs</td>
<td>int2</td>
<td></td>
<td>Number of argument strings passed to trigger function</td>
</tr>
<tr>
<td>tgattr</td>
<td>int2vector</td>
<td>pg_attribute.atanl-</td>
<td>Column numbers, if trigger is column-specific; otherwise an empty array</td>
</tr>
<tr>
<td></td>
<td></td>
<td>um</td>
<td></td>
</tr>
<tr>
<td>tgargs</td>
<td>bytea</td>
<td></td>
<td>Argument strings to pass to trigger, each NULL-terminated</td>
</tr>
<tr>
<td>tgqual</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression tree (in nodeToString() representation) for the trigger's</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>WHEN condition, or null if none</td>
</tr>
<tr>
<td>tgoldtable</td>
<td>name</td>
<td></td>
<td>REFERENCING clause name for OLD TABLE, or null if none</td>
</tr>
</tbody>
</table>
Currently, column-specific triggering is supported only for UPDATE events, and so \texttt{tgattr} is relevant only for that event type. \texttt{tgtype} might contain bits for other event types as well, but those are presumed to be table-wide regardless of what is in \texttt{tgattr}.

\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|}
\hline
\textbf{Name} & \textbf{Type} & \textbf{References} & \textbf{Description} \\
\hline
tgnewtable & name & REFERENCING clause name for NEW TABLE, or null if none & \\
\hline
\end{tabular}
\end{table}

\section*{Note}

When \texttt{tgconstraint} is nonzero, \texttt{tgconstrrelid}, \texttt{tgconstrindid}, \texttt{tgdeferrable}, and \texttt{tginitdeferred} are largely redundant with the referenced \texttt{pg_constraint} entry. However, it is possible for a non-deferrable trigger to be associated with a deferrable constraint: foreign key constraints can have some deferrable and some non-deferrable triggers.

\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|}
\hline
\textbf{Name} & \textbf{Type} & \textbf{References} & \textbf{Description} \\
\hline
oid & oid & Row identifier (hidden attribute; must be explicitly selected) & \\
\hline
cfgname & name & Text search configuration name & \\
\hline
cfgnamespace & oid & \texttt{pg_namespace.oid} & The OID of the namespace that contains this configuration \\
\hline
cfgowner & oid & \texttt{pg_authid.oid} & Owner of the configuration \\
\hline
cfgparser & oid & \texttt{pg_ts_parser.oid} & The OID of the text search parser for this configuration \\
\hline
\end{tabular}
\end{table}

\section*{52.57. \texttt{pg\_ts\_config}}

The \texttt{pg\_ts\_config} catalog contains entries representing text search configurations. A configuration specifies a particular text search parser and a list of dictionaries to use for each of the parser’s output token types. The parser is shown in the \texttt{pg\_ts\_config} entry, but the token-to-dictionary mapping is defined by subsidiary entries in \texttt{pg\_ts\_config\_map}.

PostgreSQL’s text search features are described at length in Chapter 12.

\section*{Table 52.57. \texttt{pg\_ts\_config} Columns}

\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|}
\hline
\textbf{Name} & \textbf{Type} & \textbf{References} & \textbf{Description} \\
\hline
oid & oid & Row identifier (hidden attribute; must be explicitly selected) & \\
\hline
cfgname & name & Text search configuration name & \\
\hline
cfgnamespace & oid & \texttt{pg_namespace.oid} & The OID of the namespace that contains this configuration \\
\hline
cfgowner & oid & \texttt{pg_authid.oid} & Owner of the configuration \\
\hline
cfgparser & oid & \texttt{pg_ts_parser.oid} & The OID of the text search parser for this configuration \\
\hline
\end{tabular}
\end{table}

\section*{52.58. \texttt{pg\_ts\_config\_map}}
The `pg_ts_config_map` catalog contains entries showing which text search dictionaries should be consulted, and in what order, for each output token type of each text search configuration's parser.

PostgreSQL's text search features are described at length in Chapter 12.

Table 52.58. `pg_ts_config_map` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>mapcfg</td>
<td>oid</td>
<td><code>pg_ts_config.oid</code></td>
<td>The OID of the <code>pg_ts_config</code> entry owning this map entry</td>
</tr>
<tr>
<td>maptokentype</td>
<td>integer</td>
<td></td>
<td>A token type emitted by the configuration's parser</td>
</tr>
<tr>
<td>mapseqno</td>
<td>integer</td>
<td></td>
<td>Order in which to consult this entry (lower <code>mapseqno</code>s first)</td>
</tr>
<tr>
<td>mapdict</td>
<td>oid</td>
<td><code>pg_ts_dict.oid</code></td>
<td>The OID of the text search dictionary to consult</td>
</tr>
</tbody>
</table>

52.59. `pg_ts_dict`

The `pg_ts_dict` catalog contains entries defining text search dictionaries. A dictionary depends on a text search template, which specifies all the implementation functions needed; the dictionary itself provides values for the user-settable parameters supported by the template. This division of labor allows dictionaries to be created by unprivileged users. The parameters are specified by a text string `dictinitoption`, whose format and meaning vary depending on the template.

PostgreSQL's text search features are described at length in Chapter 12.

Table 52.59. `pg_ts_dict` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>dictname</td>
<td>name</td>
<td></td>
<td>Text search dictionary name</td>
</tr>
<tr>
<td>dictnamespace</td>
<td>oid</td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace that contains this dictionary</td>
</tr>
<tr>
<td>dictowner</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>Owner of the dictionary</td>
</tr>
<tr>
<td>dicttemplate</td>
<td>oid</td>
<td><code>pg_ts_template.oid</code></td>
<td>The OID of the text search template for this dictionary</td>
</tr>
<tr>
<td>dictinitoption</td>
<td>text</td>
<td></td>
<td>Initialization option string for the template</td>
</tr>
</tbody>
</table>

52.60. `pg_ts_parser`

The `pg_ts_parser` catalog contains entries defining text search parsers. A parser is responsible for splitting input text into lexemes and assigning a token type to each lexeme. Since a parser must be implemented by C-language-level functions, creation of new parsers is restricted to database superusers.
PostgreSQL’s text search features are described at length in Chapter 12.

### Table 52.60. `pg_ts_parser` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>prsname</td>
<td>name</td>
<td></td>
<td>Text search parser name</td>
</tr>
<tr>
<td>prsnamespace</td>
<td>oid</td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace that contains this parser</td>
</tr>
<tr>
<td>prsstart</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the parser's startup function</td>
</tr>
<tr>
<td>prstoken</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the parser's next-token function</td>
</tr>
<tr>
<td>prsend</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the parser's shutdown function</td>
</tr>
<tr>
<td>prsheadline</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the parser's headline function</td>
</tr>
<tr>
<td>prslextype</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the parser's lex-type function</td>
</tr>
</tbody>
</table>

### 52.61. `pg_ts_template`

The `pg_ts_template` catalog contains entries defining text search templates. A template is the implementation skeleton for a class of text search dictionaries. Since a template must be implemented by C-language-level functions, creation of new templates is restricted to database superusers.

PostgreSQL’s text search features are described at length in Chapter 12.

### Table 52.61. `pg_ts_template` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td><code>oid</code></td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>tmplname</td>
<td><code>name</code></td>
<td></td>
<td>Text search template name</td>
</tr>
<tr>
<td>tmplnamespace</td>
<td><code>oid</code></td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace that contains this template</td>
</tr>
<tr>
<td>tmplinit</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the template’s initialization function</td>
</tr>
<tr>
<td>tmpllexize</td>
<td><code>regproc</code></td>
<td><code>pg_proc.oid</code></td>
<td>OID of the template’s lexize function</td>
</tr>
</tbody>
</table>

### 52.62. `pg_type`

The catalog `pg_type` stores information about data types. Base types and enum types (scalar types) are created with `CREATE TYPE`, and domains with `CREATE DOMAIN`. A composite type is auto-
matically created for each table in the database, to represent the row structure of the table. It is also possible to create composite types with `CREATE TYPE AS`.

**Table 52.62. pg_type Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>typname</td>
<td>name</td>
<td></td>
<td>Data type name</td>
</tr>
<tr>
<td>typnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this type</td>
</tr>
<tr>
<td>typowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the type</td>
</tr>
<tr>
<td>typlen</td>
<td>int2</td>
<td></td>
<td>For a fixed-size type, <code>typlen</code> is the number of bytes in the internal representation of the type. But for a variable-length type, <code>typlen</code> is negative. -1 indicates a “varlena” type (one that has a length word), -2 indicates a null-terminated C string.</td>
</tr>
<tr>
<td>typbyval</td>
<td>bool</td>
<td></td>
<td><code>typbyval</code> determines whether internal routines pass a value of this type by value or by reference. <code>typbyval</code> had better be false if <code>typbylen</code> is not 1, 2, or 4 (or 8 on machines where Datum is 8 bytes). Variable-length types are always passed by reference. Note that <code>typbyval</code> can be false even if the length would allow pass-by-value.</td>
</tr>
<tr>
<td>typtype</td>
<td>char</td>
<td></td>
<td><code>typtype</code> is b for a base type, c for a composite type (e.g., a table’s row type), d for a domain, e for an enum type, p for a pseudo-type, or r for a range type. See also <code>typreloid</code> and <code>typbasetype</code>.</td>
</tr>
<tr>
<td>typcategory</td>
<td>char</td>
<td></td>
<td><code>typcategory</code> is an arbitrary classification of data types that is used by the parser to</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>--------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>typispreferred</td>
<td>bool</td>
<td></td>
<td>determine which implicit casts should be &quot;preferred&quot;. See Table 52.63.</td>
</tr>
<tr>
<td>typisdefined</td>
<td>bool</td>
<td></td>
<td>True if the type is a preferred cast target within its typcategory.</td>
</tr>
<tr>
<td>typdelim</td>
<td>char</td>
<td></td>
<td>True if the type is defined, false if this is a placeholder entry for a not-yet-defined type. When typisdefined is false, nothing except the type name, namespace, and OID can be relied on.</td>
</tr>
<tr>
<td>typrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>Character that separates two values of this type when parsing array input. Note that the delimiter is associated with the array element data type, not the array data type.</td>
</tr>
<tr>
<td>typelem</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>If this is a composite type (see typtype), then this column points to the pg_class entry that defines the corresponding table. (For a free-standing composite type, the pg_class entry doesn't really represent a table, but it is needed anyway for the type's pg_attribute entries to link to.) Zero for non-composite types.</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>If typelem is not 0 then it identifies another row in pg_type. The current type can then be subscripted like an array yielding values of type typelem. A &quot;true&quot; array type is variable length (typlen = -1), but some fixed-length (typlen &gt; 0) types also have nonzero typelem, for example name and point. If a fixed-length type has a</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>typarray</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>If typarray is not 0 then it identifies another row in pg_type, which is the “true” array type having this type as element</td>
</tr>
<tr>
<td>typinput</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Input conversion function (text format)</td>
</tr>
<tr>
<td>typoutput</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Output conversion function (text format)</td>
</tr>
<tr>
<td>typreceive</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Input conversion function (binary format), or 0 if none</td>
</tr>
<tr>
<td>typsend</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Output conversion function (binary format), or 0 if none</td>
</tr>
<tr>
<td>typmodin</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Type modifier input function, or 0 if type does not support modifiers</td>
</tr>
<tr>
<td>typmodout</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Type modifier output function, or 0 to use the standard format</td>
</tr>
<tr>
<td>typanalyze</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Custom ANALYZE function, or 0 to use the standard function</td>
</tr>
<tr>
<td>typalign</td>
<td>char</td>
<td></td>
<td>typalign is the alignment required when storing a value of this type. It applies to storage on disk as well as most representations of the value inside PostgreSQL. When multiple values are stored consecutively, such as in the representation of a complete row on disk, padding is inserted before a datum of this type so that it begins on the specified boundary. The alignment reference is 2061</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------</td>
<td>----------</td>
<td>------------</td>
<td>-------------</td>
</tr>
</tbody>
</table>
| typstorage | char     |            | the beginning of the first datum in the sequence. Possible values are:  
|          |          |            | • $c = \text{char}$ alignment, i.e., no alignment needed.  
|          |          |            | • $s = \text{short}$ alignment (2 bytes on most machines).  
|          |          |            | • $i = \text{int}$ alignment (4 bytes on most machines).  
|          |          |            | • $d = \text{double}$ alignment (8 bytes on many machines, but by no means all). |

**Note**

For types used in system tables, it is critical that the size and alignment defined in `pg_type` agree with the way that the compiler will lay out the column in a structure representing a table row.

typstorage | char | typstorage tells for varlena types (those with `typlen = -1`) if
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>typnotnull</td>
<td>bool</td>
<td></td>
<td>typnotnull represents a not-null constraint on a type. Used for domains only.</td>
</tr>
<tr>
<td>typbasetype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>If this is a domain (see typtype), then typbasetype identifies the type that this one is based on. Zero if this type is not a domain.</td>
</tr>
<tr>
<td>typtypmod</td>
<td>int4</td>
<td></td>
<td>Domains use typtypmod to record the typmod to be applied to their base type (-1 if base type does not use a typmod). -1 if this type is not a domain.</td>
</tr>
<tr>
<td>typndims</td>
<td>int4</td>
<td></td>
<td>typndims is the number of array dimensions for a domain over an array (that is, typbasetype is an array type). Zero for types other</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>------------</td>
<td>------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>typcollation</td>
<td>oid</td>
<td>pg_collation.oid</td>
<td>typcollation specifies the collation of the type. If the type does not support collations, this will be zero. A base type that supports collations will have DEFAULT_COLLATION_OID here. A domain over a collatable type can have some other collation OID, if one was specified for the domain.</td>
</tr>
<tr>
<td>typdefaultbin</td>
<td>pg_node_tree</td>
<td></td>
<td>If typdefaultbin is not null, it is the nodeToString() representation of a default expression for the type. This is only used for domains.</td>
</tr>
<tr>
<td>typdefault</td>
<td>text</td>
<td></td>
<td>typdefault is null if the type has no associated default value. If typdefaultbin is not null, typdefault must contain a human-readable version of the default expression represented by typdefaultbin. If typdefaultbin is null and typdefault is not, then typdefault is the external representation of the type's default value, which can be fed to the type's input converter to produce a constant.</td>
</tr>
<tr>
<td>typacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details.</td>
</tr>
</tbody>
</table>

Table 52.63 lists the system-defined values of typcategory. Any future additions to this list will also be upper-case ASCII letters. All other ASCII characters are reserved for user-defined categories.

Table 52.63. typcategory Codes

<table>
<thead>
<tr>
<th>Code</th>
<th>Category</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>Array types</td>
</tr>
<tr>
<td>B</td>
<td>Boolean types</td>
</tr>
</tbody>
</table>
Table 52.64. pg_user_mapping Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>umuser</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>OID of the local role being mapped, 0 if the user mapping is public</td>
</tr>
<tr>
<td>umserver</td>
<td>oid</td>
<td>pg_foreign_server.oid</td>
<td>The OID of the foreign server that contains this mapping</td>
</tr>
<tr>
<td>umoptions</td>
<td>text[]</td>
<td></td>
<td>User mapping specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

52.63. pg_user_mapping

The catalog pg_user_mapping stores the mappings from local user to remote. Access to this catalog is restricted from normal users, use the view pg_user_mappings instead.

52.64. System Views

In addition to the system catalogs, PostgreSQL provides a number of built-in views. Some system views provide convenient access to some commonly used queries on the system catalogs. Other views provide access to internal server state.

The information schema (Chapter 37) provides an alternative set of views which overlap the functionality of the system views. Since the information schema is SQL-standard whereas the views described here are PostgreSQL-specific, it’s usually better to use the information schema if it provides all the information you need.

Table 52.65 lists the system views described here. More detailed documentation of each view follows below. There are some additional views that provide access to the results of the statistics collector; they are described in Table 28.2.
Except where noted, all the views described here are read-only.

**Table 52.65. System Views**

<table>
<thead>
<tr>
<th>View Name</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_available_extensions</code></td>
<td>available extensions</td>
</tr>
<tr>
<td><code>pg_available_extension_versions</code></td>
<td>available versions of extensions</td>
</tr>
<tr>
<td><code>pg_config</code></td>
<td>compile-time configuration parameters</td>
</tr>
<tr>
<td><code>pg_cursors</code></td>
<td>open cursors</td>
</tr>
<tr>
<td><code>pg_file_settings</code></td>
<td>summary of configuration file contents</td>
</tr>
<tr>
<td><code>pg_group</code></td>
<td>groups of database users</td>
</tr>
<tr>
<td><code>pg_hba_file_rules</code></td>
<td>summary of client authentication configuration file contents</td>
</tr>
<tr>
<td><code>pg_indexes</code></td>
<td>indexes</td>
</tr>
<tr>
<td><code>pg_locks</code></td>
<td>locks currently held or awaited</td>
</tr>
<tr>
<td><code>pg_matviews</code></td>
<td>materialized views</td>
</tr>
<tr>
<td><code>pg_policies</code></td>
<td>policies</td>
</tr>
<tr>
<td><code>pg_prepared_statements</code></td>
<td>prepared statements</td>
</tr>
<tr>
<td><code>pg_prepared_xacts</code></td>
<td>prepared transactions</td>
</tr>
<tr>
<td><code>pg_publication_tables</code></td>
<td>publications and their associated tables</td>
</tr>
<tr>
<td><code>pg_replication_origin_status</code></td>
<td>information about replication origins, including replication progress</td>
</tr>
<tr>
<td><code>pg_replication_slots</code></td>
<td>replication slot information</td>
</tr>
<tr>
<td><code>pg_roles</code></td>
<td>database roles</td>
</tr>
<tr>
<td><code>pg_rules</code></td>
<td>rules</td>
</tr>
<tr>
<td><code>pg_seclabels</code></td>
<td>security labels</td>
</tr>
<tr>
<td><code>pg_sequences</code></td>
<td>sequences</td>
</tr>
<tr>
<td><code>pg_settings</code></td>
<td>parameter settings</td>
</tr>
<tr>
<td><code>pg_shadow</code></td>
<td>database users</td>
</tr>
<tr>
<td><code>pg_stats</code></td>
<td>planner statistics</td>
</tr>
<tr>
<td><code>pg_tables</code></td>
<td>tables</td>
</tr>
<tr>
<td><code>pg_timezone_abbrevs</code></td>
<td>time zone abbreviations</td>
</tr>
<tr>
<td><code>pg_timezone_names</code></td>
<td>time zone names</td>
</tr>
<tr>
<td><code>pg_user</code></td>
<td>database users</td>
</tr>
<tr>
<td><code>pg_user_mappings</code></td>
<td>user mappings</td>
</tr>
<tr>
<td><code>pg_views</code></td>
<td>views</td>
</tr>
</tbody>
</table>

### 52.65. `pg_available_extensions`

The `pg_available_extensions` view lists the extensions that are available for installation. See also the `pg_extension` catalog, which shows the extensions currently installed.

**Table 52.66. `pg_available_extensions` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
<td>Extension name</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>-------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>default_version</td>
<td>text</td>
<td>Name of default version, or NULL if none is specified</td>
</tr>
<tr>
<td>installed_version</td>
<td>text</td>
<td>Currently installed version of the extension, or NULL if not installed</td>
</tr>
<tr>
<td>comment</td>
<td>text</td>
<td>Comment string from the extension's control file</td>
</tr>
</tbody>
</table>

The `pg_available_extensions` view is read only.

### 52.66. pg_available_extension_versions

The `pg_available_extension_versions` view lists the specific extension versions that are available for installation. See also the `pg_extension` catalog, which shows the extensions currently installed.

**Table 52.67. pg_available_extension_versions Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
<td>Extension name</td>
</tr>
<tr>
<td>version</td>
<td>text</td>
<td>Version name</td>
</tr>
<tr>
<td>installed</td>
<td>bool</td>
<td>True if this version of this extension is currently installed</td>
</tr>
<tr>
<td>superuser</td>
<td>bool</td>
<td>True if only superusers are allowed to install this extension</td>
</tr>
<tr>
<td>relocatable</td>
<td>bool</td>
<td>True if extension can be relocated to another schema</td>
</tr>
<tr>
<td>schema</td>
<td>name</td>
<td>Name of the schema that the extension must be installed into, or NULL if partially or fully relocatable</td>
</tr>
<tr>
<td>requires</td>
<td>name[]</td>
<td>Names of prerequisite extensions, or NULL if none</td>
</tr>
<tr>
<td>comment</td>
<td>text</td>
<td>Comment string from the extension's control file</td>
</tr>
</tbody>
</table>

The `pg_available_extension_versions` view is read only.

### 52.67. pg_config

The view `pg_config` describes the compile-time configuration parameters of the currently installed version of PostgreSQL. It is intended, for example, to be used by software packages that want to interface to PostgreSQL to facilitate finding the required header files and libraries. It provides the same basic information as the `pg_config` PostgreSQL client application.

By default, the `pg_config` view can be read only by superusers.

**Table 52.68. pg_config Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>The parameter value</td>
</tr>
</tbody>
</table>
52.68. **pg_cursors**

The `pg_cursors` view lists the cursors that are currently available. Cursors can be defined in several ways:

- via the DECLARE statement in SQL
- via the Bind message in the frontend/backend protocol, as described in Section 53.2.3
- via the Server Programming Interface (SPI), as described in Section 47.1

The `pg_cursors` view displays cursors created by any of these means. Cursors only exist for the duration of the transaction that defines them, unless they have been declared WITH HOLD. Therefore non-holdable cursors are only present in the view until the end of their creating transaction.

**Note**

Cursors are used internally to implement some of the components of PostgreSQL, such as procedural languages. Therefore, the `pg_cursors` view might include cursors that have not been explicitly created by the user.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The name of the cursor</td>
</tr>
<tr>
<td>statement</td>
<td>text</td>
<td>The verbatim query string submitted to declare this cursor</td>
</tr>
<tr>
<td>is_holdable</td>
<td>boolean</td>
<td>true if the cursor is holdable (that is, it can be accessed after the transaction that declared the cursor has committed); false otherwise</td>
</tr>
<tr>
<td>is_binary</td>
<td>boolean</td>
<td>true if the cursor was declared BINARY; false otherwise</td>
</tr>
<tr>
<td>is_scrollable</td>
<td>boolean</td>
<td>true if the cursor is scrollable (that is, it allows rows to be retrieved in a nonsequential manner); false otherwise</td>
</tr>
<tr>
<td>creation_time</td>
<td>timestamptz</td>
<td>The time at which the cursor was declared</td>
</tr>
</tbody>
</table>

The `pg_cursors` view is read only.

52.69. **pg_file_settings**

The view `pg_file_settings` provides a summary of the contents of the server's configuration file(s). A row appears in this view for each “name = value” entry appearing in the files, with annotations indicating whether the value could be applied successfully. Additional row(s) may appear for problems not linked to a “name = value” entry, such as syntax errors in the files.

This view is helpful for checking whether planned changes in the configuration files will work, or for diagnosing a previous failure. Note that this view reports on the current contents of the files, not on what was last applied by the server. (The `pg_settings` view is usually sufficient to determine that.)
By default, the `pg_file_settings` view can be read only by superusers.

Table 52.70. `pg_file_settings` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sourcefile</td>
<td>text</td>
<td>Full path name of the configuration file</td>
</tr>
<tr>
<td>sourceline</td>
<td>integer</td>
<td>Line number within the configuration file where the entry appears</td>
</tr>
<tr>
<td>seqno</td>
<td>integer</td>
<td>Order in which the entries are processed (1..n)</td>
</tr>
<tr>
<td>name</td>
<td>text</td>
<td>Configuration parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>Value to be assigned to the parameter</td>
</tr>
<tr>
<td>applied</td>
<td>boolean</td>
<td>True if the value can be applied successfully</td>
</tr>
<tr>
<td>error</td>
<td>text</td>
<td>If not null, an error message indicating why this entry could not be applied</td>
</tr>
</tbody>
</table>

If the configuration file contains syntax errors or invalid parameter names, the server will not attempt to apply any settings from it, and therefore all the `applied` fields will read as false. In such a case there will be one or more rows with non-null `error` fields indicating the problem(s). Otherwise, individual settings will be applied if possible. If an individual setting cannot be applied (e.g., invalid value, or the setting cannot be changed after server start) it will have an appropriate message in the `error` field. Another way that an entry might have `applied = false` is that it is overridden by a later entry for the same parameter name; this case is not considered an error so nothing appears in the `error` field.

See Section 19.1 for more information about the various ways to change run-time parameters.

52.70. `pg_group`

The view `pg_group` exists for backwards compatibility: it emulates a catalog that existed in PostgreSQL before version 8.1. It shows the names and members of all roles that are marked as not `rol-canlogin`, which is an approximation to the set of roles that are being used as groups.

Table 52.71. `pg_group` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>groname</td>
<td>name</td>
<td><code>pg_authid.rol-name</code></td>
<td>Name of the group</td>
</tr>
<tr>
<td>grosysid</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>ID of this group</td>
</tr>
<tr>
<td>grolist</td>
<td>oid[]</td>
<td><code>pg_authid.oid</code></td>
<td>An array containing the IDs of the roles in this group</td>
</tr>
</tbody>
</table>

52.71. `pg_hba_file_rules`

The view `pg_hba_file_rules` provides a summary of the contents of the client authentication configuration file, `pg_hba.conf`. A row appears in this view for each non-empty, non-comment line in the file, with annotations indicating whether the rule could be applied successfully.
This view can be helpful for checking whether planned changes in the authentication configuration file will work, or for diagnosing a previous failure. Note that this view reports on the current contents of the file, not on what was last loaded by the server.

By default, the `pg_hba_file_rules` view can be read only by superusers.

**Table 52.72. pg_hba_file_rules Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>line_number</td>
<td>integer</td>
<td>Line number of this rule in <code>pg_hba.conf</code></td>
</tr>
<tr>
<td>type</td>
<td>text</td>
<td>Type of connection</td>
</tr>
<tr>
<td>database</td>
<td>text[]</td>
<td>List of database name(s) to which this rule applies</td>
</tr>
<tr>
<td>user_name</td>
<td>text[]</td>
<td>List of user and group name(s) to which this rule applies</td>
</tr>
<tr>
<td>address</td>
<td>text</td>
<td>Host name or IP address, or one of <code>all</code>, <code>samehost</code>, or <code>samenet</code>, or null for local connections</td>
</tr>
<tr>
<td>netmask</td>
<td>text</td>
<td>IP address mask, or null if not applicable</td>
</tr>
<tr>
<td>auth_method</td>
<td>text</td>
<td>Authentication method</td>
</tr>
<tr>
<td>options</td>
<td>text[]</td>
<td>Options specified for authentication method, if any</td>
</tr>
<tr>
<td>error</td>
<td>text</td>
<td>If not null, an error message indicating why this line could not be processed</td>
</tr>
</tbody>
</table>

Usually, a row reflecting an incorrect entry will have values for only the `line_number` and `error` fields.

See Chapter 20 for more information about client authentication configuration.

**52.72. pg_indexes**

The view `pg_indexes` provides access to useful information about each index in the database.

**Table 52.73. pg_indexes Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td><code>pg_namespace.nspname</code></td>
<td>Name of schema containing table and index</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td><code>pg_class.relname</code></td>
<td>Name of table the index is for</td>
</tr>
<tr>
<td>indexname</td>
<td>name</td>
<td><code>pg_class.relname</code></td>
<td>Name of index</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td><code>pg_tablespace.spcname</code></td>
<td>Name of tablespace containing index (null if default for database)</td>
</tr>
<tr>
<td>indexdef</td>
<td>text</td>
<td></td>
<td>Index definition (a reconstructed <code>CREATE INDEX</code> command)</td>
</tr>
</tbody>
</table>
52.73. pg_locks

The view pg_locks provides access to information about the locks held by active processes within the database server. See Chapter 13 for more discussion of locking.

pg_locks contains one row per active lockable object, requested lock mode, and relevant process. Thus, the same lockable object might appear many times, if multiple processes are holding or waiting for locks on it. However, an object that currently has no locks on it will not appear at all.

There are several distinct types of lockable objects: whole relations (e.g., tables), individual pages of relations, individual tuples of relations, transaction IDs (both virtual and permanent IDs), and general database objects (identified by class OID and object OID, in the same way as in pg_description or pg_depend). Also, the right to extend a relation is represented as a separate lockable object. Also, “advisory” locks can be taken on numbers that have user-defined meanings.

Table 52.74. pg_locks Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>locktype</td>
<td>text</td>
<td></td>
<td>Type of the lockable object: relation, extend, page, tuple, transactionid, virtualxid, object, user-lock, or advisory</td>
</tr>
<tr>
<td>database</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>OID of the database in which the lock target exists, or zero if the target is a shared object, or null if the target is a transaction ID</td>
</tr>
<tr>
<td>relation</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the relation targeted by the lock, or null if the target is not a relation or part of a relation</td>
</tr>
<tr>
<td>page</td>
<td>integer</td>
<td></td>
<td>Page number targeted by the lock within the relation, or null if the target is not a relation page or tuple</td>
</tr>
<tr>
<td>tuple</td>
<td>smallint</td>
<td></td>
<td>Tuple number targeted by the lock within the page, or null if the target is not a tuple</td>
</tr>
<tr>
<td>virtualxid</td>
<td>text</td>
<td></td>
<td>Virtual ID of the transaction targeted by the lock, or null if the target is not a virtual transaction ID</td>
</tr>
<tr>
<td>transactionid</td>
<td>xid</td>
<td></td>
<td>ID of the transaction targeted by the lock, or null if the target is not a transaction ID</td>
</tr>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the system catalog containing the lock</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>---------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>OID of the lock target within its system catalog, or null if the target is not a general database object</td>
</tr>
<tr>
<td>objsubid</td>
<td>smallint</td>
<td></td>
<td>Column number targeted by the lock (the classid and objid refer to the table itself), or zero if the target is some other general database object, or null if the target is not a general database object</td>
</tr>
<tr>
<td>virtualtransaction</td>
<td>text</td>
<td></td>
<td>Virtual ID of the transaction that is holding or awaiting this lock</td>
</tr>
<tr>
<td>pid</td>
<td>integer</td>
<td></td>
<td>Process ID of the server process holding or awaiting this lock, or null if the lock is held by a prepared transaction</td>
</tr>
<tr>
<td>mode</td>
<td>text</td>
<td></td>
<td>Name of the lock mode held or desired by this process (see Section 13.3.1 and Section 13.2.3)</td>
</tr>
<tr>
<td>granted</td>
<td>boolean</td>
<td></td>
<td>True if lock is held, false if lock is awaited</td>
</tr>
<tr>
<td>fastpath</td>
<td>boolean</td>
<td></td>
<td>True if lock was taken via fast path, false if taken via main lock table</td>
</tr>
</tbody>
</table>

granted is true in a row representing a lock held by the indicated process. False indicates that this process is currently waiting to acquire this lock, which implies that at least one other process is holding or waiting for a conflicting lock mode on the same lockable object. The waiting process will sleep until the other lock is released (or a deadlock situation is detected). A single process can be waiting to acquire at most one lock at a time.

Throughout running a transaction, a server process holds an exclusive lock on the transaction’s virtual transaction ID. If a permanent ID is assigned to the transaction (which normally happens only if the transaction changes the state of the database), it also holds an exclusive lock on the transaction’s permanent transaction ID until it ends. When a process finds it necessary to wait specifically for another transaction to end, it does so by attempting to acquire share lock on the other transaction’s ID (either virtual or permanent ID depending on the situation). That will succeed only when the other transaction terminates and releases its locks.

Although tuples are a lockable type of object, information about row-level locks is stored on disk, not in memory, and therefore row-level locks normally do not appear in this view. If a process is waiting for a row-level lock, it will usually appear in the view as waiting for the permanent transaction ID of the current holder of that row lock.
Advisory locks can be acquired on keys consisting of either a single bigint value or two integer values. A bigint key is displayed with its high-order half in the classid column, its low-order half in the objid column, and objsubid equal to 1. The original bigint value can be reassembled with the expression (classid::bigint << 32) | objid::bigint. Integer keys are displayed with the first key in the classid column, the second key in the objid column, and objsubid equal to 2. The actual meaning of the keys is up to the user. Advisory locks are local to each database, so the database column is meaningful for an advisory lock.

pg_locks provides a global view of all locks in the database cluster, not only those relevant to the current database. Although its relation column can be joined against pg_class.oid to identify locked relations, this will only work correctly for relations in the current database (those for which the database column is either the current database’s OID or zero).

The pid column can be joined to the pid column of the pg_stat_activity view to get more information on the session holding or awaiting each lock, for example

```sql
SELECT * FROM pg_locks pl LEFT JOIN pg_stat_activity psa
ON pl.pid = psa.pid;
```

Also, if you are using prepared transactions, the virtualtransaction column can be joined to the transaction column of the pg_prepared_xacts view to get more information on prepared transactions that hold locks. (A prepared transaction can never be waiting for a lock, but it continues to hold the locks it acquired while running.) For example:

```sql
SELECT * FROM pg_locks pl LEFT JOIN pg_prepared_xacts ppx
ON pl.virtualtransaction = '-1/' || ppx.transaction;
```

While it is possible to obtain information about which processes block which other processes by joining pg_locks against itself, this is very difficult to get right in detail. Such a query would have to encode knowledge about which lock modes conflict with which others. Worse, the pg_locks view does not expose information about which processes are ahead of which others in lock wait queues, nor information about which processes are parallel workers running on behalf of which other client sessions. It is better to use the pg_blocking_pids() function (see Table 9.60) to identify which process(es) a waiting process is blocked behind.

The pg_locks view displays data from both the regular lock manager and the predicate lock manager, which are separate systems; in addition, the regular lock manager subdivides its locks into regular and fast-path locks. This data is not guaranteed to be entirely consistent. When the view is queried, data on fast-path locks (with fastpath = true) is gathered from each backend one at a time, without freezing the state of the entire lock manager, so it is possible for locks to be taken or released while information is gathered. Note, however, that these locks are known not to conflict with any other lock currently in place. After all backends have been queried for fast-path locks, the remainder of the regular lock manager is locked as a unit, and a consistent snapshot of all remaining locks is collected as an atomic action. After unlocking the regular lock manager, the predicate lock manager is similarly locked and all predicate locks are collected as an atomic action. Thus, with the exception of fast-path locks, each lock manager will deliver a consistent set of results, but as we do not lock both lock managers simultaneously, it is possible for locks to be taken or released after we interrogate the regular lock manager and before we interrogate the predicate lock manager.

Locking the regular and/or predicate lock manager could have some impact on database performance if this view is very frequently accessed. The locks are held only for the minimum amount of time necessary to obtain data from the lock managers, but this does not completely eliminate the possibility of a performance impact.

52.74. pg_matviews

The view pg_matviews provides access to useful information about each materialized view in the database.
Table 52.75. **pg_matviews** Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.n-spname</td>
<td>Name of schema containing materialized view</td>
</tr>
<tr>
<td>matviewname</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of materialized view</td>
</tr>
<tr>
<td>matviewowner</td>
<td>name</td>
<td>pg_authid.rol-name</td>
<td>Name of materialized view's owner</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td>pg_tablespace.spcname</td>
<td>Name of tablespace containing materialized view (null if default for database)</td>
</tr>
<tr>
<td>hasindexes</td>
<td>boolean</td>
<td></td>
<td>True if materialized view has (or recently had) any indexes</td>
</tr>
<tr>
<td>ispopulated</td>
<td>boolean</td>
<td></td>
<td>True if materialized view is currently populated</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>Materialized view definition (a reconstructed SELECT query)</td>
</tr>
</tbody>
</table>

52.75. **pg_policies**

The view **pg_policies** provides access to useful information about each row-level security policy in the database.

Table 52.76. **pg_policies** Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.n-spname</td>
<td>Name of schema containing table policy is on</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table policy is on</td>
</tr>
<tr>
<td>policyname</td>
<td>name</td>
<td>pg_policy.pol-name</td>
<td>Name of policy</td>
</tr>
<tr>
<td>polpermissive</td>
<td>text</td>
<td></td>
<td>Is the policy permissive or restrictive?</td>
</tr>
<tr>
<td>roles</td>
<td>name[]</td>
<td></td>
<td>The roles to which this policy applies</td>
</tr>
<tr>
<td>cmd</td>
<td>text</td>
<td></td>
<td>The command type to which the policy is applied</td>
</tr>
<tr>
<td>qual</td>
<td>text</td>
<td></td>
<td>The expression added to the security barrier qualifications for queries that this policy applies to</td>
</tr>
<tr>
<td>with_check</td>
<td>text</td>
<td></td>
<td>The expression added to the WITH</td>
</tr>
</tbody>
</table>
## 52.76. `pg_prepared_statements`

The `pg_prepared_statements` view displays all the prepared statements that are available in the current session. See PREPARE for more information about prepared statements.

`pg_prepared_statements` contains one row for each prepared statement. Rows are added to the view when a new prepared statement is created and removed when a prepared statement is released (for example, via the DEALLOCATE command).

### Table 52.77. `pg_prepared_statements` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The identifier of the prepared statement</td>
</tr>
<tr>
<td>statement</td>
<td>text</td>
<td>The query string submitted by the client to create this prepared statement. For prepared statements created via SQL, this is the <code>PREPARE</code> statement submitted by the client. For prepared statements created via the frontend/backend protocol, this is the text of the prepared statement itself.</td>
</tr>
<tr>
<td>prepare_time</td>
<td>timestamptz</td>
<td>The time at which the prepared statement was created</td>
</tr>
<tr>
<td>parameter_types</td>
<td>regtype[]</td>
<td>The expected parameter types for the prepared statement in the form of an array of <code>regtype</code>. The OID corresponding to an element of this array can be obtained by casting the <code>regtype</code> value to <code>oid</code>.</td>
</tr>
<tr>
<td>from_sql</td>
<td>boolean</td>
<td><code>true</code> if the prepared statement was created via the <code>PREPARE SQL</code> command; <code>false</code> if the statement was prepared via the frontend/backend protocol</td>
</tr>
</tbody>
</table>

The `pg_prepared_statements` view is read only.

## 52.77. `pg_prepared_xacts`

The view `pg_prepared_xacts` displays information about transactions that are currently prepared for two-phase commit (see PREPARE TRANSACTION for details).

`pg_prepared_xacts` contains one row per prepared transaction. An entry is removed when the transaction is committed or rolled back.
Table 52.78. pg_prepared_xacts Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>transaction</td>
<td>xid</td>
<td></td>
<td>Numeric transaction identifier of the prepared transaction</td>
</tr>
<tr>
<td>gid</td>
<td>text</td>
<td></td>
<td>Global transaction identifier that was assigned to the transaction</td>
</tr>
<tr>
<td>prepared</td>
<td>timestamp with time zone</td>
<td></td>
<td>Time at which the transaction was prepared for commit</td>
</tr>
<tr>
<td>owner</td>
<td>name</td>
<td>pg_authid.rolname</td>
<td>Name of the user that executed the transaction</td>
</tr>
<tr>
<td>database</td>
<td>name</td>
<td>pg_database.datname</td>
<td>Name of the database in which the transaction was executed</td>
</tr>
</tbody>
</table>

When the pg_prepared_xacts view is accessed, the internal transaction manager data structures are momentarily locked, and a copy is made for the view to display. This ensures that the view produces a consistent set of results, while not blocking normal operations longer than necessary. Nonetheless there could be some impact on database performance if this view is frequently accessed.

52.78. pg_publication_tables

The view pg_publication_tables provides information about the mapping between publications and the tables they contain. Unlike the underlying catalog pg_publication_rel, this view expands publications defined as FOR ALL TABLES, so for such publications there will be a row for each eligible table.

Table 52.79. pg_publication_tables Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pubname</td>
<td>name</td>
<td>pg_publication.pubname</td>
<td>Name of publication</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table</td>
</tr>
</tbody>
</table>

52.79. pg_replication_origin_status

The pg_replication_origin_status view contains information about how far replay for a certain origin has progressed. For more on replication origins see Chapter 50.

Table 52.80. pg_replication_origin_status Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>local_id</td>
<td>Oid</td>
<td>pg_replication_origin.oid</td>
<td>internal node identifier</td>
</tr>
<tr>
<td>external_id</td>
<td>text</td>
<td>pg_replication_origin.roname</td>
<td>external node identifier</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>remote_lsn</td>
<td>pg_lsn</td>
<td></td>
<td>The origin node’s LSN up to which data has been replicated.</td>
</tr>
<tr>
<td>local_lsn</td>
<td>pg_lsn</td>
<td></td>
<td>This node’s LSN at which remote_lsn has been replicated. Used to flush commit records before persisting data to disk when using asynchronous commits.</td>
</tr>
</tbody>
</table>

### 52.80. pg_replication_slots

The `pg_replication_slots` view provides a listing of all replication slots that currently exist on the database cluster, along with their current state.

For more on replication slots, see Section 26.2.6 and Chapter 49.

#### Table 52.81. pg_replication_slots Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>slot_name</td>
<td>name</td>
<td></td>
<td>A unique, cluster-wide identifier for the replication slot</td>
</tr>
<tr>
<td>plugin</td>
<td>name</td>
<td></td>
<td>The base name of the shared object containing the output plugin this logical slot is using, or null for physical slots.</td>
</tr>
<tr>
<td>slot_type</td>
<td>text</td>
<td></td>
<td>The slot type - physical or logical</td>
</tr>
<tr>
<td>datoid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>The OID of the database this slot is associated with, or null. Only logical slots have an associated database.</td>
</tr>
<tr>
<td>database</td>
<td>text</td>
<td>pg_database.datname</td>
<td>The name of the database this slot is associated with, or null. Only logical slots have an associated database.</td>
</tr>
<tr>
<td>temporary</td>
<td>boolean</td>
<td></td>
<td>True if this is a temporary replication slot. Temporary slots are not saved to disk and are automatically dropped on error or when the session has finished.</td>
</tr>
<tr>
<td>active</td>
<td>boolean</td>
<td></td>
<td>True if this slot is currently actively being used</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>active_pid</td>
<td>integer</td>
<td></td>
<td>The process ID of the session using this slot if the slot is currently actively being used. NULL if inactive.</td>
</tr>
<tr>
<td>xmin</td>
<td>xid</td>
<td></td>
<td>The oldest transaction that this slot needs the database to retain. VACUUM cannot remove tuples deleted by any later transaction.</td>
</tr>
<tr>
<td>catalog_xmin</td>
<td>xid</td>
<td></td>
<td>The oldest transaction affecting the system catalogs that this slot needs the database to retain. VACUUM cannot remove catalog tuples deleted by any later transaction.</td>
</tr>
<tr>
<td>restart_lsn</td>
<td>pg_lsn</td>
<td></td>
<td>The address (LSN) of oldest WAL which still might be required by the consumer of this slot and thus won’t be automatically removed during checkpoints. NULL if the LSN of this slot has never been reserved.</td>
</tr>
<tr>
<td>confirmed_flush_lsn</td>
<td>pg_lsn</td>
<td></td>
<td>The address (LSN) up to which the logical slot’s consumer has confirmed receiving data. Data older than this is not available anymore. NULL for physical slots.</td>
</tr>
</tbody>
</table>

#### 52.81. pg_roles

The view `pg_roles` provides access to information about database roles. This is simply a publicly readable view of `pg_authid` that blanks out the password field.

This view explicitly exposes the OID column of the underlying table, since that is needed to do joins to other catalogs.

**Table 52.82. pg_roles Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rolname</td>
<td>name</td>
<td></td>
<td>Role name</td>
</tr>
<tr>
<td>rolsuper</td>
<td>bool</td>
<td></td>
<td>Role has superuser privileges</td>
</tr>
<tr>
<td>rolinherit</td>
<td>bool</td>
<td></td>
<td>Role automatically inherits privileges of roles it is a member of</td>
</tr>
</tbody>
</table>

2078
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rolcreaterole</td>
<td>bool</td>
<td></td>
<td>Role can create more roles</td>
</tr>
<tr>
<td>rolcreatedb</td>
<td>bool</td>
<td></td>
<td>Role can create databases</td>
</tr>
<tr>
<td>rolcanlogin</td>
<td>bool</td>
<td></td>
<td>Role can log in. That is, this role can be given as the initial session author-</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>ization identifier</td>
</tr>
<tr>
<td>rolreplication</td>
<td>bool</td>
<td></td>
<td>Role is a replication role. A replication role can initiate replication con-</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>nections and create and drop replication slots.</td>
</tr>
<tr>
<td>rolconnlimit</td>
<td>int4</td>
<td></td>
<td>For roles that can log in, this sets maximum number of concurrent connections</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>this role can make. -1 means no limit.</td>
</tr>
<tr>
<td>rolpASSWORD</td>
<td>text</td>
<td></td>
<td>Not the password (always reads as ********)</td>
</tr>
<tr>
<td>rolvaliduntil</td>
<td>timestampz</td>
<td></td>
<td>Password expiry time (only used for password authentication); null if no ex-</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>piration</td>
</tr>
<tr>
<td>rolbypassrls</td>
<td>bool</td>
<td></td>
<td>Role bypasses every row level security policy. see Section 5.7 for more in-</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>formation.</td>
</tr>
<tr>
<td>rolconfig</td>
<td>text[]</td>
<td></td>
<td>Role-specific defaults for run-time configuration variables</td>
</tr>
<tr>
<td>oid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of role</td>
</tr>
</tbody>
</table>

#### 52.82. pg_rules

The view `pg_rules` provides access to useful information about query rewrite rules.

**Table 52.83. pg_rules Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.name</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablenamer</td>
<td>name</td>
<td>pg_class.relation</td>
<td>Name of table the rule is for</td>
</tr>
<tr>
<td>rulename</td>
<td>name</td>
<td>pg_rewrite.rule</td>
<td>Name of rule</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>Rule definition (a reconstructed creation command)</td>
</tr>
</tbody>
</table>
The `pg_rules` view excludes the ON SELECT rules of views and materialized views; those can be seen in `pg_views` and `pg_matviews`

### 52.83. pg_seclabels

The view `pg_seclabels` provides information about security labels. It is an easier-to-query version of the `pg_seclabel` catalog.

#### Table 52.84. pg_seclabels Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a security label on a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>objtype</td>
<td>text</td>
<td></td>
<td>The type of object to which this label applies, as text.</td>
</tr>
<tr>
<td>objnamespace</td>
<td>oid</td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace for this object, if applicable; otherwise NULL.</td>
</tr>
<tr>
<td>objname</td>
<td>text</td>
<td></td>
<td>The name of the object to which this label applies, as text.</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td><code>pg_seclabel.provider</code></td>
<td>The label provider associated with this label.</td>
</tr>
<tr>
<td>label</td>
<td>text</td>
<td><code>pg_seclabel.label</code></td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

### 52.84. pg_sequences

The view `pg_sequences` provides access to useful information about each sequence in the database.

#### Table 52.85. pg_sequences Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td><code>pg_namespace.nspname</code></td>
<td>Name of schema containing sequence</td>
</tr>
<tr>
<td>sequencename</td>
<td>name</td>
<td><code>pg_class.relname</code></td>
<td>Name of sequence</td>
</tr>
<tr>
<td>sequenceowner</td>
<td>name</td>
<td><code>pg_authid rolname</code></td>
<td>Name of sequence's owner</td>
</tr>
</tbody>
</table>
### System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_type</td>
<td>regtype</td>
<td>pg_type.oid</td>
<td>Data type of the sequence</td>
</tr>
<tr>
<td>start_value</td>
<td>bigint</td>
<td></td>
<td>Start value of the sequence</td>
</tr>
<tr>
<td>min_value</td>
<td>bigint</td>
<td></td>
<td>Minimum value of the sequence</td>
</tr>
<tr>
<td>max_value</td>
<td>bigint</td>
<td></td>
<td>Maximum value of the sequence</td>
</tr>
<tr>
<td>increment_by</td>
<td>bigint</td>
<td></td>
<td>Increment value of the sequence</td>
</tr>
<tr>
<td>cycle</td>
<td>boolean</td>
<td></td>
<td>Whether the sequence cycles</td>
</tr>
<tr>
<td>cache_size</td>
<td>bigint</td>
<td></td>
<td>Cache size of the sequence</td>
</tr>
<tr>
<td>last_value</td>
<td>bigint</td>
<td></td>
<td>The last sequence value written to disk. If caching is used, this value can be greater than the last value handed out from the sequence. Null if the sequence has not been read from yet. Also, if the current user does not have USAGE or SELECT privilege on the sequence, the value is null.</td>
</tr>
</tbody>
</table>

#### 52.85. pg_settings

The view `pg_settings` provides access to run-time parameters of the server. It is essentially an alternative interface to the SHOW and SET commands. It also provides access to some facts about each parameter that are not directly available from SHOW, such as minimum and maximum values.

**Table 52.86. pg_settings Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>Run-time configuration parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>Current value of the parameter</td>
</tr>
<tr>
<td>unit</td>
<td>text</td>
<td>Implicit unit of the parameter</td>
</tr>
<tr>
<td>category</td>
<td>text</td>
<td>Logical group of the parameter</td>
</tr>
<tr>
<td>short_desc</td>
<td>text</td>
<td>A brief description of the parameter</td>
</tr>
<tr>
<td>extra_desc</td>
<td>text</td>
<td>Additional, more detailed, description of the parameter</td>
</tr>
<tr>
<td>context</td>
<td>text</td>
<td>Context required to set the parameter's value (see below)</td>
</tr>
<tr>
<td>vartype</td>
<td>text</td>
<td>Parameter type (bool, enum, integer, real, or string)</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>----------</td>
<td>------------------------------------------------------------------</td>
</tr>
<tr>
<td>source</td>
<td>text</td>
<td>Source of the current parameter value</td>
</tr>
<tr>
<td>min_val</td>
<td>text</td>
<td>Minimum allowed value of the parameter (null for non-numeric values)</td>
</tr>
<tr>
<td>max_val</td>
<td>text</td>
<td>Maximum allowed value of the parameter (null for non-numeric values)</td>
</tr>
<tr>
<td>enumvals</td>
<td>text[]</td>
<td>Allowed values of an enum parameter (null for non-enum values)</td>
</tr>
<tr>
<td>boot_val</td>
<td>text</td>
<td>Parameter value assumed at server startup if the parameter is not otherwise set</td>
</tr>
<tr>
<td>reset_val</td>
<td>text</td>
<td>Value that RESET would reset the parameter to in the current session</td>
</tr>
<tr>
<td>sourcefile</td>
<td>text</td>
<td>Configuration file the current value was set in (null for values set from sources other than configuration files, or when examined by a user who is neither a superuser or a member of pg_read_all_settings); helpful when using include directives in configuration files</td>
</tr>
<tr>
<td>sourceline</td>
<td>integer</td>
<td>Line number within the configuration file the current value was set at (null for values set from sources other than configuration files, or when examined by a user who is neither a superuser or a member of pg_read_all_settings).</td>
</tr>
<tr>
<td>pending_restart</td>
<td>boolean</td>
<td>true if the value has been changed in the configuration file but needs a restart; or false otherwise.</td>
</tr>
</tbody>
</table>

There are several possible values of context. In order of decreasing difficulty of changing the setting, they are:

**internal**

These settings cannot be changed directly; they reflect internally determined values. Some of them may be adjustable by rebuilding the server with different configuration options, or by changing options supplied to initdb.

**postmaster**

These settings can only be applied when the server starts, so any change requires restarting the server. Values for these settings are typically stored in the postgresql.conf file, or passed on the command line when starting the server. Of course, settings with any of the lower context types can also be set at server start time.
sighup

Changes to these settings can be made in `postgresql.conf` without restarting the server. Send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf` and apply the changes. The postmaster will also forward the SIGHUP signal to its child processes so that they all pick up the new value.

superuser-backend

Changes to these settings can be made in `postgresql.conf` without restarting the server. They can also be set for a particular session in the connection request packet (for example, via libpq’s `PGOPTIONS` environment variable), but only if the connecting user is a superuser. However, these settings never change in a session after it is started. If you change them in `postgresql.conf`, send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf`. The new values will only affect subsequently-launched sessions.

backend

Changes to these settings can be made in `postgresql.conf` without restarting the server. They can also be set for a particular session in the connection request packet (for example, via libpq’s `PGOPTIONS` environment variable); any user can make such a change for their session. However, these settings never change in a session after it is started. If you change them in `postgresql.conf`, send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf`. The new values will only affect subsequently-launched sessions.

superuser

These settings can be set from `postgresql.conf`, or within a session via the `SET` command; but only superusers can change them via `SET`. Changes in `postgresql.conf` will affect existing sessions only if no session-local value has been established with `SET`.

user

These settings can be set from `postgresql.conf`, or within a session via the `SET` command. Any user is allowed to change their session-local value. Changes in `postgresql.conf` will affect existing sessions only if no session-local value has been established with `SET`.

See Section 19.1 for more information about the various ways to change these parameters.

The `pg_settings` view cannot be inserted into or deleted from, but it can be updated. An `UPDATE` applied to a row of `pg_settings` is equivalent to executing the `SET` command on that named parameter. The change only affects the value used by the current session. If an `UPDATE` is issued within a transaction that is later aborted, the effects of the `UPDATE` command disappear when the transaction is rolled back. Once the surrounding transaction is committed, the effects will persist until the end of the session, unless overridden by another `UPDATE` or `SET`.

### 52.86. pg_shadow

The view `pg_shadow` exists for backwards compatibility: it emulates a catalog that existed in PostgreSQL before version 8.1. It shows properties of all roles that are marked as `rolcanlogin` in `pg_authid`.

The name stems from the fact that this table should not be readable by the public since it contains passwords. `pg_user` is a publicly readable view on `pg_shadow` that blanks out the password field.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>usename</td>
<td>name</td>
<td><code>pg_authid.rolname</code></td>
<td>User name</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>ID of this user</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>---------</td>
<td>-----------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>usecreatedb</td>
<td>bool</td>
<td></td>
<td>User can create databases</td>
</tr>
<tr>
<td>usesuper</td>
<td>bool</td>
<td></td>
<td>User is a superuser</td>
</tr>
<tr>
<td>userepl</td>
<td>bool</td>
<td></td>
<td>User can initiate streaming replication and put the system in and out of backup mode.</td>
</tr>
<tr>
<td>usebypassrls</td>
<td>bool</td>
<td></td>
<td>User bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>passwd</td>
<td>text</td>
<td></td>
<td>Password (possibly encrypted); null if none. See pg_authid for details of how encrypted passwords are stored.</td>
</tr>
<tr>
<td>valuntil</td>
<td>abstime</td>
<td></td>
<td>Password expiry time (only used for password authentication)</td>
</tr>
<tr>
<td>useconfig</td>
<td>text[]</td>
<td></td>
<td>Session defaults for runtime configuration variables</td>
</tr>
</tbody>
</table>

### 52.87. pg_stats

The view pg_stats provides access to the information stored in the pg_statistic catalog. This view allows access only to rows of pg_statistic that correspond to tables the user has permission to read, and therefore it is safe to allow public read access to this view.

pg_stats is also designed to present the information in a more readable format than the underlying catalog — at the cost that its schema must be extended whenever new slot types are defined for pg_statistic.

#### Table 52.88. pg_stats Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table</td>
</tr>
<tr>
<td>attname</td>
<td>name</td>
<td>pg_attribute.atname</td>
<td>Name of the column described by this row</td>
</tr>
<tr>
<td>inherited</td>
<td>bool</td>
<td></td>
<td>If true, this row includes inheritance child columns, not just the values in the specified table</td>
</tr>
<tr>
<td>null_frac</td>
<td>real</td>
<td></td>
<td>Fraction of column entries that are null</td>
</tr>
<tr>
<td>avg_width</td>
<td>integer</td>
<td></td>
<td>Average width in bytes of column's entries</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>------------</td>
<td>------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>n_distinct</td>
<td>real</td>
<td></td>
<td>If greater than zero, the estimated number of distinct values in the column. If less than zero, the negative of the number of distinct values divided by the number of rows. (The negated form is used when <code>ANALYZE</code> believes that the number of distinct values is likely to increase as the table grows; the positive form is used when the column seems to have a fixed number of possible values.) For example, -1 indicates a unique column in which the number of distinct values is the same as the number of rows.</td>
</tr>
<tr>
<td>most_common_vals</td>
<td>anyarray</td>
<td></td>
<td>A list of the most common values in the column. (Null if no values seem to be more common than any others.)</td>
</tr>
<tr>
<td>most_common_freqs</td>
<td>real[]</td>
<td></td>
<td>A list of the frequencies of the most common values, i.e., number of occurrences of each divided by total number of rows. (Null when <code>most_common_vals</code> is.)</td>
</tr>
<tr>
<td>histogram_bounds</td>
<td>anyarray</td>
<td></td>
<td>A list of values that divide the column's values into groups of approximately equal population. The values in <code>most_common_vals</code>, if present, are omitted from this histogram calculation. (This column is null if the column data type does not have a <code>&lt;</code> operator or if the <code>most_common_vals</code> list accounts for the entire population.)</td>
</tr>
<tr>
<td>correlation</td>
<td>real</td>
<td></td>
<td>Statistical correlation between physical row</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------</td>
<td>------------</td>
<td>------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>ordering and logical</td>
<td></td>
<td></td>
<td>ordering of the column values. This ranges from -1 to +1. When the value is near -1 or +1, an index scan on the column will be estimated to be cheaper than when it is near zero, due to reduction of random access to the disk. (This column is null if the column data type does not have a &lt; operator.)</td>
</tr>
<tr>
<td>most_common_elements</td>
<td>anyarray</td>
<td></td>
<td>A list of non-null element values most often appearing within values of the column. (Null for scalar types.)</td>
</tr>
<tr>
<td>most_common_element_freqs</td>
<td>real[]</td>
<td></td>
<td>A list of the frequencies of the most common element values, i.e., the fraction of rows containing at least one instance of the given value. Two or three additional values follow the per-element frequencies; these are the minimum and maximum of the preceding per-element frequencies, and optionally the frequency of null elements. (Null when most_common_elements is.)</td>
</tr>
<tr>
<td>elem_count_histogram</td>
<td>real[]</td>
<td></td>
<td>A histogram of the counts of distinct non-null element values within the values of the column, followed by the average number of distinct non-null elements. (Null for scalar types.)</td>
</tr>
</tbody>
</table>

The maximum number of entries in the array fields can be controlled on a column-by-column basis using the ALTER TABLE SET STATISTICS command, or globally by setting the default_statistics_target run-time parameter.

52.88. pg_tables

The view pg_tables provides access to useful information about each table in the database.
Table 52.89. pg_tables Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.n-spname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table</td>
</tr>
<tr>
<td>tableowner</td>
<td>name</td>
<td>pg_authid.rolname</td>
<td>Name of table’s owner</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td>pg_tablespace.spcname</td>
<td>Name of tablespace containing table (null if default for database)</td>
</tr>
<tr>
<td>hasindexes</td>
<td>boolean</td>
<td>pg_class.relhasindex</td>
<td>True if table has (or recently had) any indexes</td>
</tr>
<tr>
<td>hasrules</td>
<td>boolean</td>
<td>pg_class.relhasrules</td>
<td>True if table has (or once had) rules</td>
</tr>
<tr>
<td>hastriggers</td>
<td>boolean</td>
<td>pg_class.relhastriggers</td>
<td>True if table has (or once had) triggers</td>
</tr>
<tr>
<td>rowsecurity</td>
<td>boolean</td>
<td>pg_class.relrowsecurity</td>
<td>True if row security is enabled on the table</td>
</tr>
</tbody>
</table>

52.89. pg_timezone_abbrevs

The view pg_timezone_abbrevs provides a list of time zone abbreviations that are currently recognized by the datetime input routines. The contents of this view change when the timezone_abbreviations run-time parameter is modified.

Table 52.90. pg_timezone_abbrevs Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>abbrev</td>
<td>text</td>
<td>Time zone abbreviation</td>
</tr>
<tr>
<td>utc_offset</td>
<td>interval</td>
<td>Offset from UTC (positive means east of Greenwich)</td>
</tr>
<tr>
<td>is_dst</td>
<td>boolean</td>
<td>True if this is a daylight-savings abbreviation</td>
</tr>
</tbody>
</table>

While most timezone abbreviations represent fixed offsets from UTC, there are some that have historically varied in value (see Section B.3 for more information). In such cases this view presents their current meaning.

52.90. pg_timezone_names

The view pg_timezone_names provides a list of time zone names that are recognized by SET TIMEZONE, along with their associated abbreviations, UTC offsets, and daylight-savings status. (Technically, PostgreSQL does not use UTC because leap seconds are not handled.) Unlike the abbreviations shown in pg_timezone_abbrevs, many of these names imply a set of daylight-savings transition date rules. Therefore, the associated information changes across local DST boundaries. The displayed information is computed based on the current value of CURRENT_TIMESTAMP.

Table 52.91. pg_timezone_names Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>Time zone name</td>
</tr>
</tbody>
</table>
## System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>abbrev</td>
<td>text</td>
<td>Time zone abbreviation</td>
</tr>
<tr>
<td>utc_offset</td>
<td>interval</td>
<td>Offset from UTC (positive means east of Greenwich)</td>
</tr>
<tr>
<td>is_dst</td>
<td>boolean</td>
<td>True if currently observing daylight savings</td>
</tr>
</tbody>
</table>

### 52.91. pg_user

The view `pg_user` provides access to information about database users. This is simply a publicly readable view of `pg_shadow` that blanks out the password field.

#### Table 52.92. pg_user Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>usename</td>
<td>name</td>
<td>User name</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>ID of this user</td>
</tr>
<tr>
<td>usecreatedb</td>
<td>bool</td>
<td>User can create databases</td>
</tr>
<tr>
<td>usesuper</td>
<td>bool</td>
<td>User is a superuser</td>
</tr>
<tr>
<td>userepl</td>
<td>bool</td>
<td>User can initiate streaming replication and put the system in and out of backup mode.</td>
</tr>
<tr>
<td>usebypassrls</td>
<td>bool</td>
<td>User bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>passwd</td>
<td>text</td>
<td>Not the password (always reads as *********)</td>
</tr>
<tr>
<td>valuntil</td>
<td>abstime</td>
<td>Password expiry time (only used for password authentication)</td>
</tr>
<tr>
<td>useconfig</td>
<td>text[]</td>
<td>Session defaults for run-time configuration variables</td>
</tr>
</tbody>
</table>

### 52.92. pg_user_mappings

The view `pg_user_mappings` provides access to information about user mappings. This is essentially a publicly readable view of `pg_user_mapping` that leaves out the options field if the user has no rights to use it.

#### Table 52.93. pg_user_mappings Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>umid</td>
<td>oid</td>
<td><code>pg_user_mapping.oid</code></td>
<td>OID of the user mapping</td>
</tr>
<tr>
<td>srvid</td>
<td>oid</td>
<td><code>pg_foreign_server.oid</code></td>
<td>The OID of the foreign server that contains this mapping</td>
</tr>
<tr>
<td>srvname</td>
<td>name</td>
<td><code>pg_foreign_server.srvname</code></td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>umuser</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>OID of the local role being mapped, 0 if the user mapping is public</td>
</tr>
</tbody>
</table>
To protect password information stored as a user mapping option, the `umoptions` column will read as null unless one of the following applies:

- current user is the user being mapped, and owns the server or holds USAGE privilege on it
- current user is the server owner and mapping is for PUBLIC
- current user is a superuser

### 52.93. pg_views

The view `pg_views` provides access to useful information about each view in the database.

**Table 52.94. pg_views Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing view</td>
</tr>
<tr>
<td>viewname</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of view</td>
</tr>
<tr>
<td>viewowner</td>
<td>name</td>
<td>pg_authid.rolname</td>
<td>Name of view's owner</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>View definition (a reconstructed SELECT query)</td>
</tr>
</tbody>
</table>
Chapter 53. Frontend/Backend Protocol

PostgreSQL uses a message-based protocol for communication between frontends and backends (clients and servers). The protocol is supported over TCP/IP and also over Unix-domain sockets. Port number 5432 has been registered with IANA as the customary TCP port number for servers supporting this protocol, but in practice any non-privileged port number can be used.

This document describes version 3.0 of the protocol, implemented in PostgreSQL 7.4 and later. For descriptions of the earlier protocol versions, see previous releases of the PostgreSQL documentation. A single server can support multiple protocol versions. The initial startup-request message tells the server which protocol version the client is attempting to use. If the major version requested by the client is not supported by the server, the connection will be rejected (for example, this would occur if the client requested protocol version 4.0, which does not exist as of this writing). If the minor version requested by the client is not supported by the server (e.g. the client requests version 3.1, but the server supports only 3.0), the server may either reject the connection or may respond with a NegotiateProtocolVersion message containing the highest minor protocol version which it supports. The client may then choose either to continue with the connection using the specified protocol version or to abort the connection.

In order to serve multiple clients efficiently, the server launches a new “backend” process for each client. In the current implementation, a new child process is created immediately after an incoming connection is detected. This is transparent to the protocol, however. For purposes of the protocol, the terms “backend” and “server” are interchangeable; likewise “frontend” and “client” are interchangeable.

53.1. Overview

The protocol has separate phases for startup and normal operation. In the startup phase, the frontend opens a connection to the server and authenticates itself to the satisfaction of the server. (This might involve a single message, or multiple messages depending on the authentication method being used.) If all goes well, the server then sends status information to the frontend, and finally enters normal operation. Except for the initial startup-request message, this part of the protocol is driven by the server.

During normal operation, the frontend sends queries and other commands to the backend, and the backend sends back query results and other responses. There are a few cases (such as NOTIFY) wherein the backend will send unsolicited messages, but for the most part this portion of a session is driven by frontend requests.

Termination of the session is normally by frontend choice, but can be forced by the backend in certain cases. In any case, when the backend closes the connection, it will roll back any open (incomplete) transaction before exiting.

Within normal operation, SQL commands can be executed through either of two sub-protocols. In the “simple query” protocol, the frontend just sends a textual query string, which is parsed and immediately executed by the backend. In the “extended query” protocol, processing of queries is separated into multiple steps: parsing, binding of parameter values, and execution. This offers flexibility and performance benefits, at the cost of extra complexity.

Normal operation has additional sub-protocols for special operations such as COPY.

53.1.1. Messaging Overview

All communication is through a stream of messages. The first byte of a message identifies the message type, and the next four bytes specify the length of the rest of the message (this length count includes itself, but not the message-type byte). The remaining contents of the message are determined by the
message type. For historical reasons, the very first message sent by the client (the startup message) has no initial message-type byte.

To avoid losing synchronization with the message stream, both servers and clients typically read an entire message into a buffer (using the byte count) before attempting to process its contents. This allows easy recovery if an error is detected while processing the contents. In extreme situations (such as not having enough memory to buffer the message), the receiver can use the byte count to determine how much input to skip before it resumes reading messages.

Conversely, both servers and clients must take care never to send an incomplete message. This is commonly done by marshaling the entire message in a buffer before beginning to send it. If a communications failure occurs partway through sending or receiving a message, the only sensible response is to abandon the connection, since there is little hope of recovering message-boundary synchronization.

53.1.2. Extended Query Overview

In the extended-query protocol, execution of SQL commands is divided into multiple steps. The state retained between steps is represented by two types of objects: prepared statements and portals. A prepared statement represents the result of parsing and semantic analysis of a textual query string. A prepared statement is not in itself ready to execute, because it might lack specific values for parameters. A portal represents a ready-to-execute or already-partially-executed statement, with any missing parameter values filled in. (For SELECT statements, a portal is equivalent to an open cursor, but we choose to use a different term since cursors don't handle non-SELECT statements.)

The overall execution cycle consists of a parse step, which creates a prepared statement from a textual query string; a bind step, which creates a portal given a prepared statement and values for any needed parameters; and an execute step that runs a portal's query. In the case of a query that returns rows (SELECT, SHOW, etc), the execute step can be told to fetch only a limited number of rows, so that multiple execute steps might be needed to complete the operation.

The backend can keep track of multiple prepared statements and portals (but note that these exist only within a session, and are never shared across sessions). Existing prepared statements and portals are referenced by names assigned when they were created. In addition, an “unnamed” prepared statement and portal exist. Although these behave largely the same as named objects, operations on them are optimized for the case of executing a query only once and then discarding it, whereas operations on named objects are optimized on the expectation of multiple uses.

53.1.3. Formats and Format Codes

Data of a particular data type might be transmitted in any of several different formats. As of PostgreSQL 7.4 the only supported formats are “text” and “binary”, but the protocol makes provision for future extensions. The desired format for any value is specified by a format code. Clients can specify a format code for each transmitted parameter value and for each column of a query result. Text has format code zero, binary has format code one, and all other format codes are reserved for future definition.

The text representation of values is whatever strings are produced and accepted by the input/output conversion functions for the particular data type. In the transmitted representation, there is no trailing null character; the frontend must add one to received values if it wants to process them as C strings. (The text format does not allow embedded nulls, by the way.)

Binary representations for integers use network byte order (most significant byte first). For other data types consult the documentation or source code to learn about the binary representation. Keep in mind that binary representations for complex data types might change across server versions; the text format is usually the more portable choice.

53.2. Message Flow

This section describes the message flow and the semantics of each message type. (Details of the exact representation of each message appear in Section 53.7.) There are several different sub-protocols de-
pending on the state of the connection: start-up, query, function call, copy, and termination. There are also special provisions for asynchronous operations (including notification responses and command cancellation), which can occur at any time after the start-up phase.

### 53.2.1. Start-up

To begin a session, a frontend opens a connection to the server and sends a startup message. This message includes the names of the user and of the database the user wants to connect to; it also identifies the particular protocol version to be used. (Optionally, the startup message can include additional settings for run-time parameters.) The server then uses this information and the contents of its configuration files (such as `pg_hba.conf`) to determine whether the connection is provisionally acceptable, and what additional authentication is required (if any).

The server then sends an appropriate authentication request message, to which the frontend must reply with an appropriate authentication response message (such as a password). For all authentication methods except GSSAPI, SSPI and SASL, there is at most one request and one response. In some methods, no response at all is needed from the frontend, and so no authentication request occurs. For GSSAPI, SSPI and SASL, multiple exchanges of packets may be needed to complete the authentication.

The authentication cycle ends with the server either rejecting the connection attempt (ErrorResponse), or sending AuthenticationOk.

The possible messages from the server in this phase are:

- **ErrorResponse**
  - The connection attempt has been rejected. The server then immediately closes the connection.

- **AuthenticationOk**
  - The authentication exchange is successfully completed.

- **AuthenticationKerberosV5**
  - The frontend must now take part in a Kerberos V5 authentication dialog (not described here, part of the Kerberos specification) with the server. If this is successful, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. This is no longer supported.

- **AuthenticationCleartextPassword**
  - The frontend must now send a PasswordMessage containing the password in clear-text form. If this is the correct password, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse.

- **AuthenticationMD5Password**
  - The frontend must now send a PasswordMessage containing the password (with user name) encrypted via MD5, then encrypted again using the 4-byte random salt specified in the AuthenticationMD5Password message. If this is the correct password, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. The actual PasswordMessage can be computed in SQL as `concat('md5', md5(concat(md5(concat(password, user-name)), random-salt)))`. (Keep in mind the `md5()` function returns its result as a hex string.)

- **AuthenticationSCMCredential**
  - This response is only possible for local Unix-domain connections on platforms that support SCM credential messages. The frontend must issue an SCM credential message and then send a sin-
single data byte. (The contents of the data byte are uninteresting; it's only used to ensure that the server waits long enough to receive the credential message.) If the credential is acceptable, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. (This message type is only issued by pre-9.1 servers. It may eventually be removed from the protocol specification.)

AuthenticationGSS

The frontend must now initiate a GSSAPI negotiation. The frontend will send a GSSResponse message with the first part of the GSSAPI data stream in response to this. If further messages are needed, the server will respond with AuthenticationGSSContinue.

AuthenticationSSPI

The frontend must now initiate a SSPI negotiation. The frontend will send a GSSResponse with the first part of the SSPI data stream in response to this. If further messages are needed, the server will respond with AuthenticationGSSContinue.

AuthenticationGSSContinue

This message contains the response data from the previous step of GSSAPI or SSPI negotiation (AuthenticationGSS, AuthenticationSSPI or a previous AuthenticationGSSContinue). If the GSSAPI or SSPI data in this message indicates more data is needed to complete the authentication, the frontend must send that data as another GSSResponse message. If GSSAPI or SSPI authentication is completed by this message, the server will next send AuthenticationOk to indicate successful authentication or ErrorResponse to indicate failure.

AuthenticationSASL

The frontend must now initiate a SASL negotiation, using one of the SASL mechanisms listed in the message. The frontend will send a SASLInitialResponse with the name of the selected mechanism, and the first part of the SASL data stream in response to this. If further messages are needed, the server will respond with AuthenticationSASLContinue. See Section 53.3 for details.

AuthenticationSASLContinue

This message contains challenge data from the previous step of SASL negotiation (AuthenticationSASL, or a previous AuthenticationSASLContinue). The frontend must respond with a SASL-Response message.

AuthenticationSASLFinal

SASL authentication has completed with additional mechanism-specific data for the client. The server will next send AuthenticationOk to indicate successful authentication, or an ErrorResponse to indicate failure. This message is sent only if the SASL mechanism specifies additional data to be sent from server to client at completion.

NegotiateProtocolVersion

The server does not support the minor protocol version requested by the client, but does support an earlier version of the protocol; this message indicates the highest supported minor version. This message will also be sent if the client requested unsupported protocol options (i.e. beginning with \_pq\_.) in the startup packet. This message will be followed by an ErrorResponse or a message indicating the success or failure of authentication.

If the frontend does not support the authentication method requested by the server, then it should immediately close the connection.

After having received AuthenticationOk, the frontend must wait for further messages from the server. In this phase a backend process is being started, and the frontend is just an interested bystander. It
Frontend/Backend Protocol

is still possible for the startup attempt to fail (ErrorResponse) or the server to decline support for the requested minor protocol version (NegotiateProtocolVersion), but in the normal case the backend will send some ParameterStatus messages, BackendKeyData, and finally ReadyForQuery.

During this phase the backend will attempt to apply any additional run-time parameter settings that were given in the startup message. If successful, these values become session defaults. An error causes ErrorResponse and exit.

The possible messages from the backend in this phase are:

BackendKeyData

This message provides secret-key data that the frontend must save if it wants to be able to issue cancel requests later. The frontend should not respond to this message, but should continue listening for a ReadyForQuery message.

ParameterStatus

This message informs the frontend about the current (initial) setting of backend parameters, such as client_encoding or DateStyle. The frontend can ignore this message, or record the settings for its future use; see Section 53.2.6 for more details. The frontend should not respond to this message, but should continue listening for a ReadyForQuery message.

ReadyForQuery

Start-up is completed. The frontend can now issue commands.

ErrorResponse

Start-up failed. The connection is closed after sending this message.

NoticeResponse

A warning message has been issued. The frontend should display the message but continue listening for ReadyForQuery or ErrorResponse.

The ReadyForQuery message is the same one that the backend will issue after each command cycle. Depending on the coding needs of the frontend, it is reasonable to consider ReadyForQuery as starting a command cycle, or to consider ReadyForQuery as ending the start-up phase and each subsequent command cycle.

53.2.2. Simple Query

A simple query cycle is initiated by the frontend sending a Query message to the backend. The message includes an SQL command (or commands) expressed as a text string. The backend then sends one or more response messages depending on the contents of the query command string, and finally a ReadyForQuery response message. ReadyForQuery informs the frontend that it can safely send a new command. (It is not actually necessary for the frontend to wait for ReadyForQuery before issuing another command, but the frontend must then take responsibility for figuring out what happens if the earlier command fails and already-issued later commands succeed.)

The possible response messages from the backend are:

CommandComplete

An SQL command completed normally.

CopyInResponse

The backend is ready to copy data from the frontend to a table; see Section 53.2.5.
CopyOutResponse

The backend is ready to copy data from a table to the frontend; see Section 53.2.5.

RowDescription

Indicates that rows are about to be returned in response to a SELECT, FETCH, etc query. The contents of this message describe the column layout of the rows. This will be followed by a DataRow message for each row being returned to the frontend.

DataRow

One of the set of rows returned by a SELECT, FETCH, etc query.

EmptyQueryResponse

An empty query string was recognized.

ErrorResponse

An error has occurred.

ReadyForQuery

Processing of the query string is complete. A separate message is sent to indicate this because the query string might contain multiple SQL commands. (CommandComplete marks the end of processing one SQL command, not the whole string.) ReadyForQuery will always be sent, whether processing terminates successfully or with an error.

NoticeResponse

A warning message has been issued in relation to the query. Notices are in addition to other responses, i.e., the backend will continue processing the command.

The response to a SELECT query (or other queries that return row sets, such as EXPLAIN or SHOW) normally consists of RowDescription, zero or more DataRow messages, and then CommandComplete. COPY to or from the frontend invokes special protocol as described in Section 53.2.5. All other query types normally produce only a CommandComplete message.

Since a query string could contain several queries (separated by semicolons), there might be several such response sequences before the backend finishes processing the query string. ReadyForQuery is issued when the entire string has been processed and the backend is ready to accept a new query string.

If a completely empty (no contents other than whitespace) query string is received, the response is EmptyQueryResponse followed by ReadyForQuery.

In the event of an error, ErrorResponse is issued followed by ReadyForQuery. All further processing of the query string is aborted by ErrorResponse (even if more queries remained in it). Note that this might occur partway through the sequence of messages generated by an individual query.

In simple Query mode, the format of retrieved values is always text, except when the given command is a FETCH from a cursor declared with the BINARY option. In that case, the retrieved values are in binary format. The format codes given in the RowDescription message tell which format is being used.

A frontend must be prepared to accept ErrorResponse and NoticeResponse messages whenever it is expecting any other type of message. See also Section 53.2.6 concerning messages that the backend might generate due to outside events.

Recommended practice is to code frontends in a state-machine style that will accept any message type at any time that it could make sense, rather than wiring in assumptions about the exact sequence of messages.
53.2.2.1. Multiple Statements in a Simple Query

When a simple Query message contains more than one SQL statement (separated by semicolons), those statements are executed as a single transaction, unless explicit transaction control commands are included to force a different behavior. For example, if the message contains

```
INSERT INTO mytable VALUES(1);
SELECT 1/0;
INSERT INTO mytable VALUES(2);
```

then the divide-by-zero failure in the `SELECT` will force rollback of the first `INSERT`. Furthermore, because execution of the message is abandoned at the first error, the second `INSERT` is never attempted at all.

If instead the message contains

```
BEGIN;
INSERT INTO mytable VALUES(1);
COMMIT;
INSERT INTO mytable VALUES(2);
SELECT 1/0;
```

then the first `INSERT` is committed by the explicit `COMMIT` command. The second `INSERT` and the `SELECT` are still treated as a single transaction, so that the divide-by-zero failure will roll back the second `INSERT`, but not the first one.

This behavior is implemented by running the statements in a multi-statement Query message in an implicit transaction block unless there is some explicit transaction block for them to run in. The main difference between an implicit transaction block and a regular one is that an implicit block is closed automatically at the end of the Query message, either by an implicit commit if there was no error, or an implicit rollback if there was an error. This is similar to the implicit commit or rollback that happens for a statement executed by itself (when not in a transaction block).

If the session is already in a transaction block, as a result of a `BEGIN` in some previous message, then the Query message simply continues that transaction block, whether the message contains one statement or several. However, if the Query message contains a `COMMIT` or `ROLLBACK` closing the existing transaction block, then any following statements are executed in an implicit transaction block. Conversely, if a `BEGIN` appears in a multi-statement Query message, then it starts a regular transaction block that will only be terminated by an explicit `COMMIT` or `ROLLBACK`, whether that appears in this Query message or a later one. If the `BEGIN` follows some statements that were executed as an implicit transaction block, those statements are not immediately committed; in effect, they are retroactively included into the new regular transaction block.

A `COMMIT` or `ROLLBACK` appearing in an implicit transaction block is executed as normal, closing the implicit block; however, a warning will be issued since a `COMMIT` or `ROLLBACK` without a previous `BEGIN` might represent a mistake. If more statements follow, a new implicit transaction block will be started for them.

Savepoints are not allowed in an implicit transaction block, since they would conflict with the behavior of automatically closing the block upon any error.

Remember that, regardless of any transaction control commands that may be present, execution of the Query message stops at the first error. Thus for example given

```
BEGIN;
SELECT 1/0;
ROLLBACK;
```
in a single Query message, the session will be left inside a failed regular transaction block, since the ROLLBACK is not reached after the divide-by-zero error. Another ROLLBACK will be needed to restore the session to a usable state.

Another behavior of note is that initial lexical and syntactic analysis is done on the entire query string before any of it is executed. Thus simple errors (such as a misspelled keyword) in later statements can prevent execution of any of the statements. This is normally invisible to users since the statements would all roll back anyway when done as an implicit transaction block. However, it can be visible when attempting to do multiple transactions within a multi-statement Query. For instance, if a typo turned our previous example into

BEGIN;
INSERT INTO mytable VALUES(1);
COMMIT;
INSERT INTO mytable VALUES(2);
SELECT 1/0;

then none of the statements would get run, resulting in the visible difference that the first INSERT is not committed. Errors detected at semantic analysis or later, such as a misspelled table or column name, do not have this effect.

### 53.2.3. Extended Query

The extended query protocol breaks down the above-described simple query protocol into multiple steps. The results of preparatory steps can be re-used multiple times for improved efficiency. Furthermore, additional features are available, such as the possibility of supplying data values as separate parameters instead of having to insert them directly into a query string.

In the extended protocol, the frontend first sends a Parse message, which contains a textual query string, optionally some information about data types of parameter placeholders, and the name of a destination prepared-statement object (an empty string selects the unnamed prepared statement). The response is either ParseComplete or ErrorResponse. Parameter data types can be specified by OID; if not given, the parser attempts to infer the data types in the same way as it would do for untyped literal string constants.

---

**Note**

A parameter data type can be left unspecified by setting it to zero, or by making the array of parameter type OIDs shorter than the number of parameter symbols (5$n$) used in the query string. Another special case is that a parameter's type can be specified as void (that is, the OID of the void pseudo-type). This is meant to allow parameter symbols to be used for function parameters that are actually OUT parameters. Ordinarily there is no context in which a void parameter could be used, but if such a parameter symbol appears in a function's parameter list, it is effectively ignored. For example, a function call such as `foo($1,$2,$3,$4)` could match a function with two IN and two OUT arguments, if $3$ and $4$ are specified as having type void.

---

**Note**

The query string contained in a Parse message cannot include more than one SQL statement; else a syntax error is reported. This restriction does not exist in the simple-query protocol, but it does exist in the extended protocol, because allowing prepared statements or portals to contain multiple commands would complicate the protocol unduly.
If successfully created, a named prepared-statement object lasts till the end of the current session, unless explicitly destroyed. An unnamed prepared statement lasts only until the next Parse statement specifying the unnamed statement as destination is issued. (Note that a simple Query message also destroys the unnamed statement.) Named prepared statements must be explicitly closed before they can be redefined by another Parse message, but this is not required for the unnamed statement. Named prepared statements can also be created and accessed at the SQL command level, using PREPARE and EXECUTE.

Once a prepared statement exists, it can be readied for execution using a Bind message. The Bind message gives the name of the source prepared statement (empty string denotes the unnamed prepared statement), the name of the destination portal (empty string denotes the unnamed portal), and the values to use for any parameter placeholders present in the prepared statement. The supplied parameter set must match those needed by the prepared statement. (If you declared any void parameters in the Parse message, pass NULL values for them in the Bind message.) Bind also specifies the format to use for any data returned by the query; the format can be specified overall, or per-column. The response is either BindComplete or ErrorResponse.

**Note**
The choice between text and binary output is determined by the format codes given in Bind, regardless of the SQL command involved. The BINARY attribute in cursor declarations is irrelevant when using extended query protocol.

Query planning typically occurs when the Bind message is processed. If the prepared statement has no parameters, or is executed repeatedly, the server might save the created plan and re-use it during subsequent Bind messages for the same prepared statement. However, it will do so only if it finds that a generic plan can be created that is not much less efficient than a plan that depends on the specific parameter values supplied. This happens transparently so far as the protocol is concerned.

If successfully created, a named portal object lasts till the end of the current transaction, unless explicitly destroyed. An unnamed portal is destroyed at the end of the transaction, or as soon as the next Bind statement specifying the unnamed portal as destination is issued. (Note that a simple Query message also destroys the unnamed portal.) Named portals must be explicitly closed before they can be redefined by another Bind message, but this is not required for the unnamed portal. Named portals can also be created and accessed at the SQL command level, using DECLARE CURSOR and FETCH.

Once a portal exists, it can be executed using an Execute message. The Execute message specifies the portal name (empty string denotes the unnamed portal) and a maximum result-row count (zero meaning "fetch all rows"). The result-row count is only meaningful for portals containing commands that return row sets; in other cases the command is always executed to completion, and the row count is ignored. The possible responses to Execute are the same as those described above for queries issued via simple query protocol, except that Execute doesn't cause ReadyForQuery or RowDescription to be issued.

If Execute terminates before completing the execution of a portal (due to reaching a nonzero result-row count), it will send a PortalSuspended message; the appearance of this message tells the frontend that another Execute should be issued against the same portal to complete the operation. The CommandComplete message indicating completion of the source SQL command is not sent until the portal's execution is completed. Therefore, an Execute phase is always terminated by the appearance of exactly one of these messages: CommandComplete, EmptyQueryResponse (if the portal was created from an empty query string), ErrorResponse, or PortalSuspended.

At completion of each series of extended-query messages, the frontend should issue a Sync message. This parameterless message causes the backend to close the current transaction if it's not inside a BEGIN/COMMIT transaction block ("close" meaning to commit if no error, or roll back if error). Then a ReadyForQuery response is issued. The purpose of Sync is to provide a resynchronization point for error recovery. When an error is detected while processing any extended-query message, the
backend issues ErrorResponse, then reads and discards messages until a Sync is reached, then issues ReadyForQuery and returns to normal message processing. (But note that no skipping occurs if an error is detected while processing Sync — this ensures that there is one and only one ReadyForQuery sent for each Sync.)

**Note**

Sync does not cause a transaction block opened with BEGIN to be closed. It is possible to detect this situation since the ReadyForQuery message includes transaction status information.

In addition to these fundamental, required operations, there are several optional operations that can be used with extended-query protocol.

The Describe message (portal variant) specifies the name of an existing portal (or an empty string for the unnamed portal). The response is a RowDescription message describing the rows that will be returned by executing the portal; or a NoData message if the portal does not contain a query that will return rows; or ErrorResponse if there is no such portal.

The Describe message (statement variant) specifies the name of an existing prepared statement (or an empty string for the unnamed prepared statement). The response is a ParameterDescription message describing the parameters needed by the statement, followed by a RowDescription message describing the rows that will be returned when the statement is eventually executed (or a NoData message if the statement will not return rows). ErrorResponse is issued if there is no such prepared statement. Note that since Bind has not yet been issued, the formats to be used for returned columns are not yet known to the backend; the format code fields in the RowDescription message will be zeroes in this case.

**Tip**

In most scenarios the frontend should issue one or the other variant of Describe before issuing Execute, to ensure that it knows how to interpret the results it will get back.

The Close message closes an existing prepared statement or portal and releases resources. It is not an error to issue Close against a nonexistent statement or portal name. The response is normally CloseComplete, but could be ErrorResponse if some difficulty is encountered while releasing resources. Note that closing a prepared statement implicitly closes any open portals that were constructed from that statement.

The Flush message does not cause any specific output to be generated, but forces the backend to deliver any data pending in its output buffers. A Flush must be sent after any extended-query command except Sync, if the frontend wishes to examine the results of that command before issuing more commands. Without Flush, messages returned by the backend will be combined into the minimum possible number of packets to minimize network overhead.

**Note**

The simple Query message is approximately equivalent to the series Parse, Bind, portal Describe, Execute, Close, Sync, using the unnamed prepared statement and portal objects and no parameters. One difference is that it will accept multiple SQL statements in the query string, automatically performing the bind/describe/execute sequence for each one in succession. Another difference is that it will not return ParseComplete, BindComplete, CloseComplete, or NoData messages.
53.2.4. Function Call

The Function Call sub-protocol allows the client to request a direct call of any function that exists in the database's pg_proc system catalog. The client must have execute permission for the function.

Note

The Function Call sub-protocol is a legacy feature that is probably best avoided in new code. Similar results can be accomplished by setting up a prepared statement that does SELECT function($1, ...). The Function Call cycle can then be replaced with Bind/Execute.

A Function Call cycle is initiated by the frontend sending a FunctionCall message to the backend. The backend then sends one or more response messages depending on the results of the function call, and finally a ReadyForQuery response message. ReadyForQuery informs the frontend that it can safely send a new query or function call.

The possible response messages from the backend are:

ErrorResponse

An error has occurred.

FunctionCallResponse

The function call was completed and returned the result given in the message. (Note that the Function Call protocol can only handle a single scalar result, not a row type or set of results.)

ReadyForQuery

Processing of the function call is complete. ReadyForQuery will always be sent, whether processing terminates successfully or with an error.

NoticeResponse

A warning message has been issued in relation to the function call. Notices are in addition to other responses, i.e., the backend will continue processing the command.

53.2.5. COPY Operations

The COPY command allows high-speed bulk data transfer to or from the server. Copy-in and copy-out operations each switch the connection into a distinct sub-protocol, which lasts until the operation is completed.

Copy-in mode (data transfer to the server) is initiated when the backend executes a COPY FROM STDIN SQL statement. The backend sends a CopyInResponse message to the frontend. The frontend should then send zero or more CopyData messages, forming a stream of input data. (The message boundaries are not required to have anything to do with row boundaries, although that is often a reasonable choice.) The frontend can terminate the copy-in mode by sending either a CopyDone message (allowing successful termination) or a CopyFail message (which will cause the COPY SQL statement to fail with an error). The backend then reverts to the command-processing mode it was in before the COPY started, which will be either simple or extended query protocol. It will next send either CommandComplete (if successful) or ErrorResponse (if not).

In the event of a backend-detected error during copy-in mode (including receipt of a CopyFail message), the backend will issue an ErrorResponse message. If the COPY command was issued via an extended-query message, the backend will now discard frontend messages until a Sync message is received, then it will issue ReadyForQuery and return to normal processing. If the COPY command was
issued in a simple Query message, the rest of that message is discarded and ReadyForQuery is issued. In either case, any subsequent CopyData, CopyDone, or CopyFail messages issued by the frontend will simply be dropped.

The backend will ignore Flush and Sync messages received during copy-in mode. Receipt of any other non-copy message type constitutes an error that will abort the copy-in state as described above. (The exception for Flush and Sync is for the convenience of client libraries that always send Flush or Sync after an Execute message, without checking whether the command to be executed is a COPY FROM STDIN.)

Copy-out mode (data transfer from the server) is initiated when the backend executes a COPY TO STDOUT SQL statement. The backend sends a CopyOutResponse message to the frontend, followed by zero or more CopyData messages (always one per row), followed by CopyDone. The backend then reverts to the command-processing mode it was in before the COPY started, and sends CommandComplete. The frontend cannot abort the transfer (except by closing the connection or issuing a Cancel request), but it can discard unwanted CopyData and CopyDone messages.

In the event of a backend-detected error during copy-out mode, the backend will issue an ErrorResponse message and revert to normal processing. The frontend should treat receipt of ErrorResponse as terminating the copy-out mode.

It is possible for NoticeResponse and ParameterStatus messages to be interspersed between CopyData messages; frontends must handle these cases, and should be prepared for other asynchronous message types as well (see Section 53.2.6). Otherwise, any message type other than CopyData or CopyDone may be treated as terminating copy-out mode.

There is another Copy-related mode called copy-both, which allows high-speed bulk data transfer to and from the server. Copy-both mode is initiated when a backend in walsender mode executes a START_REPLICATION statement. The backend sends a CopyBothResponse message to the frontend. Both the backend and the frontend may then send CopyData messages until either end sends a CopyDone message. After the client sends a CopyDone message, the connection goes from copy-both mode to copy-out mode, and the client may not send any more CopyData messages. Similarly, when the server sends a CopyDone message, the connection goes into copy-in mode, and the server may not send any more CopyData messages. After both sides have sent a CopyDone message, the copy mode is terminated, and the backend reverts to the command-processing mode. In the event of a backend-detected error during copy-both mode, the backend will issue an ErrorResponse message, discard frontend messages until a Sync message is received, and then issue ReadyForQuery and return to normal processing. The frontend should treat receipt of ErrorResponse as terminating the copy in both directions; no CopyDone should be sent in this case. See Section 53.4 for more information on the subprotocol transmitted over copy-both mode.

The CopyInResponse, CopyOutResponse and CopyBothResponse messages include fields that inform the frontend of the number of columns per row and the format codes being used for each column. (As of the present implementation, all columns in a given COPY operation will use the same format, but the message design does not assume this.)

53.2.6. Asynchronous Operations

There are several cases in which the backend will send messages that are not specifically prompted by the frontend’s command stream. Frontends must be prepared to deal with these messages at any time, even when not engaged in a query. At minimum, one should check for these cases before beginning to read a query response.

It is possible for NoticeResponse messages to be generated due to outside activity; for example, if the database administrator commands a “fast” database shutdown, the backend will send a NoticeResponse indicating this fact before closing the connection. Accordingly, frontends should always be prepared to accept and display NoticeResponse messages, even when the connection is nominally idle.

ParameterStatus messages will be generated whenever the active value changes for any of the parameters the backend believes the frontend should know about. Most commonly this occurs in response to a
SET SQL command executed by the frontend, and this case is effectively synchronous — but it is also possible for parameter status changes to occur because the administrator changed a configuration file and then sent the SIGHUP signal to the server. Also, if a SET command is rolled back, an appropriate ParameterStatus message will be generated to report the current effective value.

At present there is a hard-wired set of parameters for which ParameterStatus will be generated: they are server_version, server_encoding, client_encoding, application_name, is_superuser, session_authorization, DateStyle, IntervalStyle, TimeZone, integer_datetimes, and standard_conforming_strings. (server_encoding, TimeZone, and integer_datetimes were not reported by releases before 8.0; standard_conforming_strings was not reported by releases before 8.1; application_name was not reported by releases before 9.0.) Note that server_version, server_encoding and integer_datetimes are pseudo-parameters that cannot change after startup. This set might change in the future, or even become configurable. Accordingly, a frontend should simply ignore ParameterStatus for parameters that it does not understand or care about.

If a frontend issues a LISTEN command, then the backend will send a NotificationResponse message (not to be confused with NoticeResponse!) whenever a NOTIFY command is executed for the same channel name.

### Note

At present, NotificationResponse can only be sent outside a transaction, and thus it will not occur in the middle of a command-response series, though it might occur just before ReadyForQuery. It is unwise to design frontend logic that assumes that, however. Good practice is to be able to accept NotificationResponse at any point in the protocol.

#### 53.2.7. Canceling Requests in Progress

During the processing of a query, the frontend might request cancellation of the query. The cancel request is not sent directly on the open connection to the backend for reasons of implementation efficiency: we don't want to have the backend constantly checking for new input from the frontend during query processing. Cancel requests should be relatively infrequent, so we make them slightly cumbersome in order to avoid a penalty in the normal case.

To issue a cancel request, the frontend opens a new connection to the server and sends a CancelRequest message, rather than the StartupMessage message that would ordinarily be sent across a new connection. The server will process this request and then close the connection. For security reasons, no direct reply is made to the cancel request message.

A CancelRequest message will be ignored unless it contains the same key data (PID and secret key) passed to the frontend during connection start-up. If the request matches the PID and secret key for a currently executing backend, the processing of the current query is aborted. (In the existing implementation, this is done by sending a special signal to the backend process that is processing the query.)

The cancellation signal might or might not have any effect — for example, if it arrives after the backend has finished processing the query, then it will have no effect. If the cancellation is effective, it results in the current command being terminated early with an error message.

The upshot of all this is that for reasons of both security and efficiency, the frontend has no direct way to tell whether a cancel request has succeeded. It must continue to wait for the backend to respond to the query. Issuing a cancel simply improves the odds that the current query will finish soon, and improves the odds that it will fail with an error message instead of succeeding.

Since the cancel request is sent across a new connection to the server and not across the regular frontend/backend communication link, it is possible for the cancel request to be issued by any process,
not just the frontend whose query is to be canceled. This might provide additional flexibility when building multiple-process applications. It also introduces a security risk, in that unauthorized persons might try to cancel queries. The security risk is addressed by requiring a dynamically generated secret key to be supplied in cancel requests.

### 53.2.8. Termination

The normal, graceful termination procedure is that the frontend sends a Terminate message and immediately closes the connection. On receipt of this message, the backend closes the connection and terminates.

In rare cases (such as an administrator-commanded database shutdown) the backend might disconnect without any frontend request to do so. In such cases the backend will attempt to send an error or notice message giving the reason for the disconnection before it closes the connection.

Other termination scenarios arise from various failure cases, such as core dump at one end or the other, loss of the communications link, loss of message-boundary synchronization, etc. If either frontend or backend sees an unexpected closure of the connection, it should clean up and terminate. The frontend has the option of launching a new backend by recontacting the server if it doesn't want to terminate itself. Closing the connection is also advisable if an unrecognizable message type is received, since this probably indicates loss of message-boundary sync.

For either normal or abnormal termination, any open transaction is rolled back, not committed. One should note however that if a frontend disconnects while a non-SELECT query is being processed, the backend will probably finish the query before noticing the disconnection. If the query is outside any transaction block (BEGIN ... COMMIT sequence) then its results might be committed before the disconnection is recognized.

### 53.2.9. SSL Session Encryption

If PostgreSQL was built with SSL support, frontend/backend communications can be encrypted using SSL. This provides communication security in environments where attackers might be able to capture the session traffic. For more information on encrypting PostgreSQL sessions with SSL, see Section 18.9.

To initiate an SSL-encrypted connection, the frontend initially sends an SSLRequest message rather than a StartupMessage. The server then responds with a single byte containing $S$ or $N$, indicating that it is willing or unwilling to perform SSL, respectively. The frontend might close the connection at this point if it is dissatisfied with the response. To continue after $S$, perform an SSL startup handshake (not described here, part of the SSL specification) with the server. If this is successful, continue with sending the usual StartupMessage. In this case the StartupMessage and all subsequent data will be SSL-encrypted. To continue after $N$, send the usual StartupMessage and proceed without encryption.

The frontend should also be prepared to handle an ErrorMessage response to SSLRequest from the server. This would only occur if the server predates the addition of SSL support to PostgreSQL. (Such servers are now very ancient, and likely do not exist in the wild anymore.) In this case the connection must be closed, but the frontend might choose to open a fresh connection and proceed without requesting SSL.

An initial SSLRequest can also be used in a connection that is being opened to send a CancelRequest message.

While the protocol itself does not provide a way for the server to force SSL encryption, the administrator can configure the server to reject unencrypted sessions as a byproduct of authentication checking.

### 53.3. SASL Authentication

SASL is a framework for authentication in connection-oriented protocols. At the moment, PostgreSQL implements two SASL authentication mechanisms, SCRAM-SHA-256 and SCRAM-SHA-256-
PLUS. More might be added in the future. The below steps illustrate how SASL authentication is performed in general, while the next subsection gives more details on SCRAM-SHA-256 and SCRAM-SHA-256-PLUS.

**SASL Authentication Message Flow**

1. To begin a SASL authentication exchange, the server sends an AuthenticationSASL message. It includes a list of SASL authentication mechanisms that the server can accept, in the server's preferred order.

2. The client selects one of the supported mechanisms from the list, and sends a SASLInitialResponse message to the server. The message includes the name of the selected mechanism, and an optional Initial Client Response, if the selected mechanism uses that.

3. One or more server-challenge and client-response message will follow. Each server-challenge is sent in an AuthenticationSASLContinue message, followed by a response from client in an SASLResponse message. The particulars of the messages are mechanism specific.

4. Finally, when the authentication exchange is completed successfully, the server sends an AuthenticationSASLFinal message, followed immediately by an AuthenticationOk message. The AuthenticationSASLFinal contains additional server-to-client data, whose content is particular to the selected authentication mechanism. If the authentication mechanism doesn't use additional data that's sent at completion, the AuthenticationSASLFinal message is not sent.

On error, the server can abort the authentication at any stage, and send an ErrorMessage.

### 53.3.1. SCRAM-SHA-256 authentication

The implemented SASL mechanisms at the moment are SCRAM-SHA-256 and its variant with channel binding SCRAM-SHA-256-PLUS. They are described in detail in RFC 7677 and RFC 5802.

When SCRAM-SHA-256 is used in PostgreSQL, the server will ignore the user name that the client sends in the client-first-message. The user name that was already sent in the startup message is used instead. PostgreSQL supports multiple character encodings, while SCRAM dictates UTF-8 to be used for the user name, so it might be impossible to represent the PostgreSQL user name in UTF-8.

The SCRAM specification dictates that the password is also in UTF-8, and is processed with the **SASLprep** algorithm. PostgreSQL, however, does not require UTF-8 to be used for the password. When a user's password is set, it is processed with SASLprep as if it was in UTF-8, regardless of the actual encoding used. However, if it is not a legal UTF-8 byte sequence, or it contains UTF-8 byte sequences that are prohibited by the SASLprep algorithm, the raw password will be used without SASLprep processing, instead of throwing an error. This allows the password to be normalized when it is in UTF-8, but still allows a non-UTF-8 password to be used, and doesn't require the system to know which encoding the password is in.

**Channel binding** is supported in PostgreSQL builds with SSL support. The SASL mechanism name for SCRAM with channel binding is SCRAM-SHA-256-PLUS. The channel binding type used by PostgreSQL is tls-server-end-point.

In SCRAM without channel binding, the server chooses a random number that is transmitted to the client to be mixed with the user-supplied password in the transmitted password hash. While this prevents the password hash from being successfully retransmitted in a later session, it does not prevent a fake server between the real server and client from passing through the server's random value and successfully authenticating.

SCRAM with channel binding prevents such man-in-the-middle attacks by mixing the signature of the server's certificate into the transmitted password hash. While a fake server can retransmit the real server's certificate, it doesn't have access to the private key matching that certificate, and therefore cannot prove it is the owner, causing SSL connection failure.
Example

1. The server sends an AuthenticationSASL message. It includes a list of SASL authentication mechanisms that the server can accept. This will be SCRAM-SHA-256-PLUS and SCRAM-SHA-256 if the server is built with SSL support, or else just the latter.

2. The client responds by sending a SASLInitialResponse message, which indicates the chosen mechanism, SCRAM-SHA-256 or SCRAM-SHA-256-PLUS. (A client is free to choose either mechanism, but for better security it should choose the channel-binding variant if it can support it.) In the Initial Client response field, the message contains the SCRAM client-first-message. The client-first-message also contains the channel binding type chosen by the client.

3. Server sends an AuthenticationSASLContinue message, with a SCRAM server-first message as the content.

4. Client sends a SASLResponse message, with SCRAM client-final-message as the content.

5. Server sends an AuthenticationSASLFinal message, with the SCRAM server-final-message, followed immediately by an AuthenticationOk message.

53.4. Streaming Replication Protocol

To initiate streaming replication, the frontend sends the replication parameter in the startup message. A Boolean value of true (or on, yes, 1) tells the backend to go into physical replication walsender mode, wherein a small set of replication commands, shown below, can be issued instead of SQL statements.

Passing database as the value for the replication parameter instructs the backend to go into logical replication walsender mode, connecting to the database specified in the dbname parameter. In logical replication walsender mode, the replication commands shown below as well as normal SQL commands can be issued.

In either physical replication or logical replication walsender mode, only the simple query protocol can be used.

For the purpose of testing replication commands, you can make a replication connection via psql or any other libpq-using tool with a connection string including the replication option, e.g.:

```
psql "dbname=postgres replication=database" -c "IDENTIFY_SYSTEM;"
```

However, it is often more useful to use pg_receivewal (for physical replication) or pg_recvlogical (for logical replication).

Replication commands are logged in the server log when log_replication_commands is enabled.

The commands accepted in replication mode are:

**IDENTIFY_SYSTEM**

Requests the server to identify itself. Server replies with a result set of a single row, containing four fields:

```
systemid(text)
```

The unique system identifier identifying the cluster. This can be used to check that the base backup used to initialize the standby came from the same cluster.
timeline (int4)

Current timeline ID. Also useful to check that the standby is consistent with the master.

xlogpos (text)

Current WAL flush location. Useful to get a known location in the write-ahead log where
streaming can start.

dbname (text)

Database connected to or null.

SHOW name

Requests the server to send the current setting of a run-time parameter. This is similar to the SQL
command SHOW.

name

The name of a run-time parameter. Available parameters are documented in Chapter 19.

TIMELINE_HISTORY tli

Requests the server to send over the timeline history file for timeline tli. Server replies with a
result set of a single row, containing two fields:

filename (text)

File name of the timeline history file, e.g., 00000002.history.

content (bytea)

Contents of the timeline history file.

CREATE_REPLICATION_SLOT slot_name [ TEMPORARY ] [ PHYSICAL [ RESERVE_WAL ] ]
LOGICAL output_plugin [ EXPORT_SNAPSHOT | NOEXPORT_SNAPSHOT | USE_SNAPSHOT ]

Create a physical or logical replication slot. See Section 26.2.6 for more about replication slots.

slot_name

The name of the slot to create. Must be a valid replication slot name (see Section 26.2.6.1).

output_plugin

The name of the output plugin used for logical decoding (see Section 49.6).

TEMPORARY

Specify that this replication slot is a temporary one. Temporary slots are not saved to disk
and are automatically dropped on error or when the session has finished.

RESERVE_WAL

Specify that this physical replication slot reserves WAL immediately. Otherwise, WAL is
only reserved upon connection from a streaming replication client.

EXPORT_SNAPSHOT
NOEXPORT_SNAPSHOT
USE_SNAPSHOT

Decides what to do with the snapshot created during logical slot initialization. EXPORT_SNAP-
SHOTS, which is the default, will export the snapshot for use in other sessions. This op-
tion can't be used inside a transaction. USE_SNAPSHOT will use the snapshot for the current transaction executing the command. This option must be used in a transaction, and CREATE_REPLICATION_SLOT must be the first command run in that transaction. Finally, NOEXPORT_SNAPSHOT will just use the snapshot for logical decoding as normal but won't do anything else with it.

In response to this command, the server will send a one-row result set containing the following fields:

- **slot_name (text)**
  
  The name of the newly-created replication slot.

- **consistent_point (text)**
  
  The WAL location at which the slot became consistent. This is the earliest location from which streaming can start on this replication slot.

- **snapshot_name (text)**
  
  The identifier of the snapshot exported by the command. The snapshot is valid until a new command is executed on this connection or the replication connection is closed. Null if the created slot is physical.

- **output_plugin (text)**
  
  The name of the output plugin used by the newly-created replication slot. Null if the created slot is physical.

**START_REPLICATION [SLOT slot_name] [PHYSICAL] XXX/XXX [TIMELINE tli]**

Instructs server to start streaming WAL, starting at WAL location XXX/XXX. If TIMELINE option is specified, streaming starts on timeline tli; otherwise, the server's current timeline is selected. The server can reply with an error, for example if the requested section of WAL has already been recycled. On success, server responds with a CopyBothResponse message, and then starts to stream WAL to the frontend.

If a slot's name is provided via slot_name, it will be updated as replication progresses so that the server knows which WAL segments, and if hot_standby_feedback is on which transactions, are still needed by the standby.

If the client requests a timeline that's not the latest but is part of the history of the server, the server will stream all the WAL on that timeline starting from the requested start point up to the point where the server switched to another timeline. If the client requests streaming at exactly the end of an old timeline, the server responds immediately with CommandComplete without entering COPY mode.

After streaming all the WAL on a timeline that is not the latest one, the server will end streaming by exiting the COPY mode. When the client acknowledges this by also exiting COPY mode, the server sends a result set with one row and two columns, indicating the next timeline in this server's history. The first column is the next timeline's ID (type int8), and the second column is the WAL location where the switch happened (type text). Usually, the switch position is the end of the WAL that was streamed, but there are corner cases where the server can send some WAL from the old timeline that it has not itself replayed before promoting. Finally, the server sends CommandComplete message, and is ready to accept a new command.

WAL data is sent as a series of CopyData messages. (This allows other information to be intermixed; in particular the server can send an ErrorResponse message if it encounters a failure after beginning to stream.) The payload of each CopyData message from server to the client contains a message of one of the following formats:
XLogData (B)

Byte1(‘w’)

Identifies the message as WAL data.

Int64

The starting point of the WAL data in this message.

Int64

The current end of WAL on the server.

Int64

The server’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte1

A section of the WAL data stream.

A single WAL record is never split across two XLogData messages. When a WAL record crosses a WAL page boundary, and is therefore already split using continuation records, it can be split at the page boundary. In other words, the first main WAL record and its continuation records can be sent in different XLogData messages.

Primary keepalive message (B)

Byte1(‘k’)

Identifies the message as a sender keepalive.

Int64

The current end of WAL on the server.

Int64

The server’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte1

1 means that the client should reply to this message as soon as possible, to avoid a timeout disconnect. 0 otherwise.

The receiving process can send replies back to the sender at any time, using one of the following message formats (also in the payload of a CopyData message):

Standby status update (F)

Byte1(‘r’)

Identifies the message as a receiver status update.

Int64

The location of the last WAL byte + 1 received and written to disk in the standby.

Int64

The location of the last WAL byte + 1 flushed to disk in the standby.
Int64

The location of the last WAL byte + 1 applied in the standby.

Int64

The client's system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte1

If 1, the client requests the server to reply to this message immediately. This can be used to ping the server, to test if the connection is still healthy.

Hot Standby feedback message (F)

Byte1('h')

Identifies the message as a Hot Standby feedback message.

Int64

The client's system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Int32

The standby's current global xmin, excluding the catalog_xmin from any replication slots. If both this value and the following catalog_xmin are 0 this is treated as a notification that Hot Standby feedback will no longer be sent on this connection. Later non-zero messages may reinitiate the feedback mechanism.

Int32

The epoch of the global xmin xid on the standby.

Int32

The lowest catalog_xmin of any replication slots on the standby. Set to 0 if no catalog_xmin exists on the standby or if hot standby feedback is being disabled.

Int32

The epoch of the catalog_xmin xid on the standby.

START_REPLICATION SLOT slot_name LOGICAL XXX/XXX [ ( option_name [ option_value] [, ...] ) ]

Instructs server to start streaming WAL for logical replication, starting at WAL location XXX/XXX. The server can reply with an error, for example if the requested section of WAL has already been recycled. On success, server responds with a CopyBothResponse message, and then starts to stream WAL to the frontend.

The messages inside the CopyBothResponse messages are of the same format documented for START_REPLICATION ... PHYSICAL.

The output plugin associated with the selected slot is used to process the output for streaming.

SLOT slot_name

The name of the slot to stream changes from. This parameter is required, and must correspond to an existing logical replication slot created with CREATE_REPLICATION_SLOT in LOGICAL mode.
The WAL location to begin streaming at.

The name of an option passed to the slot's logical decoding plugin.

Optional value, in the form of a string constant, associated with the specified option.

Drops a replication slot, freeing any reserved server-side resources. If the slot is a logical slot that was created in a database other than the database the walsender is connected to, this command fails.

The name of the slot to drop.

This option causes the command to wait if the slot is active until it becomes inactive, instead of the default behavior of raising an error.

Instructs the server to start streaming a base backup. The system will automatically be put in backup mode before the backup is started, and taken out of it when the backup is complete. The following options are accepted:

Sets the label of the backup. If none is specified, a backup label of `base backup` will be used. The quoting rules for the label are the same as a standard SQL string with standard_conforming_strings turned on.

Request information required to generate a progress report. This will send back an approximate size in the header of each tablespace, which can be used to calculate how far along the stream is done. This is calculated by enumerating all the file sizes once before the transfer is even started, and might as such have a negative impact on the performance. In particular, it might take longer before the first data is streamed. Since the database files can change during the backup, the size is only approximate and might both grow and shrink between the time of approximation and the sending of the actual files.

Request a fast checkpoint.

Include the necessary WAL segments in the backup. This will include all the files between start and stop backup in the `pg_wal` directory of the base directory tar file.

By default, the backup will wait until the last required WAL segment has been archived, or emit a warning if log archiving is not enabled. Specifying `NOWAIT` disables both the waiting and the warning, leaving the client responsible for ensuring the required log is available.
MAX_RATE rate

Limit (throttle) the maximum amount of data transferred from server to client per unit of time. The expected unit is kilobytes per second. If this option is specified, the value must either be equal to zero or it must fall within the range from 32 kB through 1 GB (inclusive). If zero is passed or the option is not specified, no restriction is imposed on the transfer.

TABLESPACE_MAP

Include information about symbolic links present in the directory pg_tblspc in a file named tablespace_map. The tablespace map file includes each symbolic link name as it exists in the directory pg_tblspc/ and the full path of that symbolic link.

NOVERIFY_CHECKSUMS

By default, checksums are verified during a base backup if they are enabled. Specifying NOVERIFY_CHECKSUMS disables this verification.

When the backup is started, the server will first send two ordinary result sets, followed by one or more CopyResponse results.

The first ordinary result set contains the starting position of the backup, in a single row with two columns. The first column contains the start position given in XLogRecPtr format, and the second column contains the corresponding timeline ID.

The second ordinary result set has one row for each tablespace. The fields in this row are:

spcoid(oid)

The OID of the tablespace, or null if it's the base directory.

spclocation(text)

The full path of the tablespace directory, or null if it's the base directory.

size(int8)

The approximate size of the tablespace, if progress report has been requested; otherwise it's null.

After the second regular result set, one or more CopyResponse results will be sent, one for the main data directory and one for each additional tablespace other than pg_default and pg_global. The data in the CopyResponse results will be a tar format (following the “ustar interchange format” specified in the POSIX 1003.1-2008 standard) dump of the tablespace contents, except that the two trailing blocks of zeroes specified in the standard are omitted. After the tar data is complete, a final ordinary result set will be sent, containing the WAL end position of the backup, in the same format as the start position.

The tar archive for the data directory and each tablespace will contain all files in the directories, regardless of whether they are PostgreSQL files or other files added to the same directory. The only excluded files are:

• postmaster.pid
• postmaster.opts
• pg_internal.init (found in multiple directories)
• Various temporary files and directories created during the operation of the PostgreSQL server, such as any file or directory beginning with pgsql_tmp and temporary relations.
• Unlogged relations, except for the init fork which is required to recreate the (empty) unlogged relation on recovery.
• pg_wal, including subdirectories. If the backup is run with WAL files included, a synthesized version of pg_wal will be included, but it will only contain the files necessary for the backup to work, not the rest of the contents.
53.5. Logical Streaming Replication Protocol

This section describes the logical replication protocol, which is the message flow started by the START_REPLICATION SLOT slot_name LOGICAL replication command. The logical streaming replication protocol builds on the primitives of the physical streaming replication protocol.

53.5.1. Logical Streaming Replication Parameters

The logical replication START_REPLICATION command accepts following parameters:

- `proto_version`
  Protocol version. Currently only version 1 is supported.

- `publication_names`
  Comma separated list of publication names for which to subscribe (receive changes). The individual publication names are treated as standard objects names and can be quoted the same as needed.

53.5.2. Logical Replication Protocol Messages

The individual protocol messages are discussed in the following subsections. Individual messages are described in Section 53.9.

All top-level protocol messages begin with a message type byte. While represented in code as a character, this is a signed byte with no associated encoding.

Since the streaming replication protocol supplies a message length there is no need for top-level protocol messages to embed a length in their header.

53.5.3. Logical Replication Protocol Message Flow

With the exception of the START_REPLICATION command and the replay progress messages, all information flows only from the backend to the frontend.

The logical replication protocol sends individual transactions one by one. This means that all messages between a pair of Begin and Commit messages belong to the same transaction.

Every sent transaction contains zero or more DML messages (Insert, Update, Delete). In case of a cascaded setup it can also contain Origin messages. The origin message indicated that the transaction originated on different replication node. Since a replication node in the scope of logical replication protocol can be pretty much anything, the only identifier is the origin name. It's downstream's responsibility to handle this as needed (if needed). The Origin message is always sent before any DML messages in the transaction.

Every DML message contains an arbitrary relation ID, which can be mapped to an ID in the Relation messages. The Relation messages describe the schema of the given relation. The Relation message is...
sent for a given relation either because it is the first time we send a DML message for given relation
in the current session or because the relation definition has changed since the last Relation message
was sent for it. The protocol assumes that the client is capable of caching the metadata for as many
relations as needed.

53.6. Message Data Types

This section describes the base data types used in messages.

Int\(n\)\(\hat{\alpha}\)

An \(n\)-bit integer in network byte order (most significant byte first). If \(\hat{\alpha}\) is specified it is the exact
value that will appear, otherwise the value is variable. Eg. Int16, Int32(42).

Int\(n\)[\(k\)]

An array of \(k\) \(n\)-bit integers, each in network byte order. The array length \(k\) is always determined
by an earlier field in the message. Eg. Int16[M].

String\(s\)

A null-terminated string (C-style string). There is no specific length limitation on strings. If \(s\) is
specified it is the exact value that will appear, otherwise the value is variable. Eg. String,
String("user").

Note

There is no predefined limit on the length of a string that can be returned by the
backend. Good coding strategy for a frontend is to use an expandable buffer so
that anything that fits in memory can be accepted. If that's not feasible, read the
full string and discard trailing characters that don't fit into your fixed-size buffer.

Byte\(n\)(\(c\))

Exactly \(n\) bytes. If the field width \(n\) is not a constant, it is always determinable from an earlier
field in the message. If \(c\) is specified it is the exact value. Eg. Byte2, Byte1('n').

53.7. Message Formats

This section describes the detailed format of each message. Each is marked to indicate that it can be
sent by a frontend (F), a backend (B), or both (F & B). Notice that although each message includes a
byte count at the beginning, the message format is defined so that the message end can be found without
reference to the byte count. This aids validity checking. (The CopyData message is an exception,
because it forms part of a data stream; the contents of any individual CopyData message cannot be
interpretable on their own.)

AuthenticationOk (B)

Byte1('R')

Identifies the message as an authentication request.

Int32(8)

Length of message contents in bytes, including self.
Int32(0)
   Specifies that the authentication was successful.

AuthenticationKerberosV5 (B)
   Byte1('R')
   Identifies the message as an authentication request.
   Int32(8)
   Length of message contents in bytes, including self.
   Int32(2)
   Specifies that Kerberos V5 authentication is required.

AuthenticationCleartextPassword (B)
   Byte1('R')
   Identifies the message as an authentication request.
   Int32(8)
   Length of message contents in bytes, including self.
   Int32(3)
   Specifies that a clear-text password is required.

AuthenticationMD5Password (B)
   Byte1('R')
   Identifies the message as an authentication request.
   Int32(12)
   Length of message contents in bytes, including self.
   Int32(5)
   Specifies that an MD5-encrypted password is required.
   Byte4
   The salt to use when encrypting the password.

AuthenticationSCMCredential (B)
   Byte1('R')
   Identifies the message as an authentication request.
   Int32(8)
   Length of message contents in bytes, including self.
   Int32(6)
   Specifies that an SCM credentials message is required.
AuthenticationGSS (B)
Byte1('R')
Identifies the message as an authentication request.
Int32(8)
Length of message contents in bytes, including self.
Int32(7)
Specifies that GSSAPI authentication is required.

AuthenticationSSPI (B)
Byte1('R')
Identifies the message as an authentication request.
Int32(8)
Length of message contents in bytes, including self.
Int32(9)
Specifies that SSPI authentication is required.

AuthenticationGSSContinue (B)
Byte1('R')
Identifies the message as an authentication request.
Int32
Length of message contents in bytes, including self.
Int32(8)
Specifies that this message contains GSSAPI or SSPI data.

Byte
GSSAPI or SSPI authentication data.

AuthenticationSASL (B)
Byte1('R')
Identifies the message as an authentication request.
Int32
Length of message contents in bytes, including self.
Int32(10)
Specifies that SASL authentication is required.

The message body is a list of SASL authentication mechanisms, in the server's order of preference. A zero byte is required as terminator after the last authentication mechanism name. For each mechanism, there is the following:
String
Name of a SASL authentication mechanism.

AuthenticationSASLContinue (B)
Byte1('R')
Identifies the message as an authentication request.

Int32
Length of message contents in bytes, including self.

Int32(11)
Specifies that this message contains a SASL challenge.

Byte
SASL data, specific to the SASL mechanism being used.

AuthenticationSASLFinal (B)
Byte1('R')
Identifies the message as an authentication request.

Int32
Length of message contents in bytes, including self.

Int32(12)
Specifies that SASL authentication has completed.

Byte
SASL outcome "additional data", specific to the SASL mechanism being used.

BackendKeyData (B)
Byte1('K')
Identifies the message as cancellation key data. The frontend must save these values if it wishes to be able to issue CancelRequest messages later.

Int32(12)
Length of message contents in bytes, including self.

Int32
The process ID of this backend.

Int32
The secret key of this backend.

Bind (F)
Byte1('B')
Identifies the message as a Bind command.
Int32
Length of message contents in bytes, including self.

String
The name of the destination portal (an empty string selects the unnamed portal).

String
The name of the source prepared statement (an empty string selects the unnamed prepared statement).

Int16
The number of parameter format codes that follow (denoted C below). This can be zero to indicate that there are no parameters or that the parameters all use the default format (text); or one, in which case the specified format code is applied to all parameters; or it can equal the actual number of parameters.

Int16[C]
The parameter format codes. Each must presently be zero (text) or one (binary).

Int16
The number of parameter values that follow (possibly zero). This must match the number of parameters needed by the query.

Next, the following pair of fields appear for each parameter:

Int32
The length of the parameter value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL parameter value. No value bytes follow in the NULL case.

Byte n
The value of the parameter, in the format indicated by the associated format code. n is the above length.

After the last parameter, the following fields appear:

Int16
The number of result-column format codes that follow (denoted R below). This can be zero to indicate that there are no result columns or that the result columns should all use the default format (text); or one, in which case the specified format code is applied to all result columns (if any); or it can equal the actual number of result columns of the query.

Int16[R]
The result-column format codes. Each must presently be zero (text) or one (binary).

BindComplete (B)

Byte1('2')
Identifies the message as a Bind-complete indicator.
Int32(4)

Length of message contents in bytes, including self.

CancelRequest (F)

Int32(16)

Length of message contents in bytes, including self.

Int32(80877102)

The cancel request code. The value is chosen to contain 1234 in the most significant 16 bits, and 5678 in the least significant 16 bits. (To avoid confusion, this code must not be the same as any protocol version number.)

Int32

The process ID of the target backend.

Int32

The secret key for the target backend.

Close (F)

Byte1('C')

Identifies the message as a Close command.

Int32

Length of message contents in bytes, including self.

Byte1

'S' to close a prepared statement; or 'P' to close a portal.

String

The name of the prepared statement or portal to close (an empty string selects the unnamed prepared statement or portal).

CloseComplete (B)

Byte1('3')

Identifies the message as a Close-complete indicator.

Int32(4)

Length of message contents in bytes, including self.

CommandComplete (B)

Byte1('C')

Identifies the message as a command-completed response.

Int32

Length of message contents in bytes, including self.
String

The command tag. This is usually a single word that identifies which SQL command was
completed.

For an INSERT command, the tag is INSERT oid rows, where rows is the number of
rows inserted. oid is the object ID of the inserted row if rows is 1 and the target table has
OIDs; otherwise oid is 0.

For a DELETE command, the tag is DELETE rows where rows is the number of rows
deleted.

For an UPDATE command, the tag is UPDATE rows where rows is the number of rows
updated.

For a SELECT or CREATE TABLE AS command, the tag is SELECT rows where rows
is the number of rows retrieved.

For a MOVE command, the tag is MOVE rows where rows is the number of rows the cursor's
position has been changed by.

For a FETCH command, the tag is FETCH rows where rows is the number of rows that
have been retrieved from the cursor.

For a COPY command, the tag is COPY rows where rows is the number of rows copied.
(Note: the row count appears only in PostgreSQL 8.2 and later.)

CopyData (F & B)

Byte1('d')

Identifies the message as COPY data.

Int32

Length of message contents in bytes, including self.

Byte'n'

Data that forms part of a COPY data stream. Messages sent from the backend will always
correspond to single data rows, but messages sent by frontends might divide the data stream
arbitrarily.

CopyDone (F & B)

Byte1('c')

Identifies the message as a COPY-complete indicator.

Int32(4)

Length of message contents in bytes, including self.

CopyFail (F)

Byte1('f')

Identifies the message as a COPY-failure indicator.

Int32

Length of message contents in bytes, including self.
String

An error message to report as the cause of failure.

CopyInResponse (B)

Byte1('G')

Identifies the message as a Start Copy In response. The frontend must now send copy-in data (if not prepared to do so, send a CopyFail message).

Int32

Length of message contents in bytes, including self.

Int8

0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16

The number of columns in the data to be copied (denoted $N$ below).

Int16[$N$]

The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

CopyOutResponse (B)

Byte1('H')

Identifies the message as a Start Copy Out response. This message will be followed by copy-out data.

Int32

Length of message contents in bytes, including self.

Int8

0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16

The number of columns in the data to be copied (denoted $N$ below).

Int16[$N$]

The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

CopyBothResponse (B)

Byte1('W')

Identifies the message as a Start Copy Both response. This message is used only for Streaming Replication.
Int32
Length of message contents in bytes, including self.

Int8
0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16
The number of columns in the data to be copied (denoted N below).

Int16[N]
The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

DataRow (B)
Byte1('D')
Identifies the message as a data row.

Int32
Length of message contents in bytes, including self.

Int16
The number of column values that follow (possibly zero).

Next, the following pair of fields appear for each column:

Int32
The length of the column value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL column value. No value bytes follow in the NULL case.

Byte n
The value of the column, in the format indicated by the associated format code. n is the above length.

Describe (F)
Byte1('D')
Identifies the message as a Describe command.

Int32
Length of message contents in bytes, including self.

Byte1
'S' to describe a prepared statement; or 'P' to describe a portal.

String
The name of the prepared statement or portal to describe (an empty string selects the unnamed prepared statement or portal).
EmptyQueryResponse (B)

Byte1('I')

Identifies the message as a response to an empty query string. (This substitutes for CommandComplete.)

Int32(4)

Length of message contents in bytes, including self.

ErrorResponse (B)

Byte1('E')

Identifies the message as an error.

Int32

Length of message contents in bytes, including self.

The message body consists of one or more identified fields, followed by a zero byte as a terminator. Fields can appear in any order. For each field there is the following:

Byte1

A code identifying the field type; if zero, this is the message terminator and no string follows. The presently defined field types are listed in Section 53.8. Since more field types might be added in future, frontends should silently ignore fields of unrecognized type.

String

The field value.

Execute (F)

Byte1('E')

Identifies the message as an Execute command.

Int32

Length of message contents in bytes, including self.

String

The name of the portal to execute (an empty string selects the unnamed portal).

Int32

Maximum number of rows to return, if portal contains a query that returns rows (ignored otherwise). Zero denotes “no limit”.

Flush (F)

Byte1('H')

Identifies the message as a Flush command.

Int32(4)

Length of message contents in bytes, including self.
FunctionCall (F)

Byte1('F')

Identifies the message as a function call.

Int32

Length of message contents in bytes, including self.

Int32

Specifies the object ID of the function to call.

Int16

The number of argument format codes that follow (denoted $C$ below). This can be zero to indicate that there are no arguments or that the arguments all use the default format (text); or one, in which case the specified format code is applied to all arguments; or it can equal the actual number of arguments.

Int16[$C$]

The argument format codes. Each must presently be zero (text) or one (binary).

Int16

Specifies the number of arguments being supplied to the function.

Next, the following pair of fields appear for each argument:

Int32

The length of the argument value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL argument value. No value bytes follow in the NULL case.

Byte$n$

The value of the argument, in the format indicated by the associated format code. $n$ is the above length.

After the last argument, the following field appears:

Int16

The format code for the function result. Must presently be zero (text) or one (binary).

FunctionCallResponse (B)

Byte1('V')

Identifies the message as a function call result.

Int32

Length of message contents in bytes, including self.

Int32

The length of the function result value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL function result. No value bytes follow in the NULL case.
Byte $n$

The value of the function result, in the format indicated by the associated format code. $n$ is the above length.

GSSResponse (F)

Byte1('p')

Identifies the message as a GSSAPI or SSPI response. Note that this is also used for SASL and password response messages. The exact message type can be deduced from the context.

Int32

Length of message contents in bytes, including self.

Byte $n$

GSSAPI/SSPI specific message data.

NegotiateProtocolVersion (B)

Byte1('v')

Identifies the message as a protocol version negotiation message.

Int32

Length of message contents in bytes, including self.

Int32

Newest minor protocol version supported by the server for the major protocol version requested by the client.

Int32

Number of protocol options not recognized by the server.

Then, for protocol option not recognized by the server, there is the following:

String

The option name.

NoData (B)

Byte1('n')

Identifies the message as a no-data indicator.

Int32(4)

Length of message contents in bytes, including self.

NoticeResponse (B)

Byte1('N')

Identifies the message as a notice.

Int32

Length of message contents in bytes, including self.
The message body consists of one or more identified fields, followed by a zero byte as a terminator. Fields can appear in any order. For each field there is the following:

**Byte1**

A code identifying the field type; if zero, this is the message terminator and no string follows. The presently defined field types are listed in Section 53.8. Since more field types might be added in future, frontends should silently ignore fields of unrecognized type.

**String**

The field value.

### NotificationResponse (B)

**Byte1('A')**

Identifies the message as a notification response.

**Int32**

Length of message contents in bytes, including self.

**Int32**

The process ID of the notifying backend process.

**String**

The name of the channel that the notify has been raised on.

**String**

The “payload” string passed from the notifying process.

### ParameterDescription (B)

**Byte1('t')**

Identifies the message as a parameter description.

**Int32**

Length of message contents in bytes, including self.

**Int16**

The number of parameters used by the statement (can be zero).

Then, for each parameter, there is the following:

**Int32**

Specify the object ID of the parameter data type.

### ParameterStatus (B)

**Byte1('S')**

Identifies the message as a run-time parameter status report.

**Int32**

Length of message contents in bytes, including self.
String
  The name of the run-time parameter being reported.

String
  The current value of the parameter.

Parse (F)

Byte1('P')
  Identifies the message as a Parse command.

Int32
  Length of message contents in bytes, including self.

String
  The name of the destination prepared statement (an empty string selects the unnamed prepared statement).

String
  The query string to be parsed.

Int16
  The number of parameter data types specified (can be zero). Note that this is not an indication of the number of parameters that might appear in the query string, only the number that the frontend wants to prespecify types for.

Then, for each parameter, there is the following:

Int32
  Specifies the object ID of the parameter data type. Placing a zero here is equivalent to leaving the type unspecified.

ParseComplete (B)

Byte1('1')
  Identifies the message as a Parse-complete indicator.

Int32(4)
  Length of message contents in bytes, including self.

PasswordMessage (F)

Byte1('p')
  Identifies the message as a password response. Note that this is also used for GSSAPI, SSPI and SASL response messages. The exact message type can be deduced from the context.

Int32
  Length of message contents in bytes, including self.

String
  The password (encrypted, if requested).
PortalSuspended (B)

Byte1('s')

Identifies the message as a portal-suspended indicator. Note this only appears if an Execute
message’s row-count limit was reached.

Int32(4)

Length of message contents in bytes, including self.

Query (F)

Byte1('Q')

Identifies the message as a simple query.

Int32

Length of message contents in bytes, including self.

String

The query string itself.

ReadyForQuery (B)

Byte1('Z')

Identifies the message type. ReadyForQuery is sent whenever the backend is ready for a new
query cycle.

Int32(5)

Length of message contents in bytes, including self.

Byte1

Current backend transaction status indicator. Possible values are 'I' if idle (not in a transaction
block); 'T' if in a transaction block; or 'E' if in a failed transaction block (queries will be
rejected until block is ended).

RowDescription (B)

Byte1('T')

Identifies the message as a row description.

Int32

Length of message contents in bytes, including self.

Int16

Specifies the number of fields in a row (can be zero).

Then, for each field, there is the following:

String

The field name.
Int32

If the field can be identified as a column of a specific table, the object ID of the table; otherwise zero.

Int16

If the field can be identified as a column of a specific table, the attribute number of the column; otherwise zero.

Int32

The object ID of the field's data type.

Int16

The data type size (see \texttt{pg\_type.typlen}). Note that negative values denote variable-width types.

Int32

The type modifier (see \texttt{pg\_attribute.atttypmod}). The meaning of the modifier is type-specific.

Int16

The format code being used for the field. Currently will be zero (text) or one (binary). In a RowDescription returned from the statement variant of Describe, the format code is not yet known and will always be zero.

\textbf{SASLInitialResponse (F)}

\texttt{Byte1('p')}

Identifies the message as an initial SASL response. Note that this is also used for GSSAPI, SSPI and password response messages. The exact message type is deduced from the context.

Int32

Length of message contents in bytes, including self.

\textbf{String}

Name of the SASL authentication mechanism that the client selected.

Int32

Length of SASL mechanism specific "Initial Client Response" that follows, or -1 if there is no Initial Response.

\texttt{Byte n}

SASL mechanism specific "Initial Response".

\textbf{SASLResponse (F)}

\texttt{Byte1('p')}

Identifies the message as a SASL response. Note that this is also used for GSSAPI, SSPI and password response messages. The exact message type can be deduced from the context.

Int32

Length of message contents in bytes, including self.
SASL mechanism specific message data.

SSLRequest (F)

Int32(8)

Length of message contents in bytes, including self.

Int32(80877103)

The SSL request code. The value is chosen to contain 1234 in the most significant 16 bits, and 5679 in the least significant 16 bits. (To avoid confusion, this code must not be the same as any protocol version number.)

StartupMessage (F)

Int32

Length of message contents in bytes, including self.

Int32(196608)

The protocol version number. The most significant 16 bits are the major version number (3 for the protocol described here). The least significant 16 bits are the minor version number (0 for the protocol described here).

The protocol version number is followed by one or more pairs of parameter name and value strings. A zero byte is required as a terminator after the last name/value pair. Parameters can appear in any order. user is required, others are optional. Each parameter is specified as:

String

The parameter name. Currently recognized names are:

user

The database user name to connect as. Required; there is no default.

database

The database to connect to. Defaults to the user name.

options

Command-line arguments for the backend. (This is deprecated in favor of setting individual run-time parameters.) Spaces within this string are considered to separate arguments, unless escaped with a backslash (\); write \ to represent a literal backslash.

replication

Used to connect in streaming replication mode, where a small set of replication commands can be issued instead of SQL statements. Value can be true, false, or database, and the default is false. See Section 53.4 for details.

In addition to the above, other parameters may be listed. Parameter names beginning with _pq_ are reserved for use as protocol extensions, while others are treated as run-time parameters to be set at backend start time. Such settings will be applied during backend start (after parsing the command-line arguments if any) and will act as session defaults.

String

The parameter value.
Sync (F)

Byte1('S')

Identifies the message as a Sync command.

Int32(4)

Length of message contents in bytes, including self.

Terminate (F)

Byte1('X')

Identifies the message as a termination.

Int32(4)

Length of message contents in bytes, including self.

53.8. Error and Notice Message Fields

This section describes the fields that can appear in ErrorResponse and NoticeResponse messages. Each field type has a single-byte identification token. Note that any given field type should appear at most once per message.

S

Severity: the field contents are ERROR, FATAL, or PANIC (in an error message), or WARNING, NOTICE, DEBUG, INFO, or LOG (in a notice message), or a localized translation of one of these. Always present.

V

Severity: the field contents are ERROR, FATAL, or PANIC (in an error message), or WARNING, NOTICE, DEBUG, INFO, or LOG (in a notice message). This is identical to the S field except that the contents are never localized. This is present only in messages generated by PostgreSQL versions 9.6 and later.

C

Code: the SQLSTATE code for the error (see Appendix A). Not localizable. Always present.

M

Message: the primary human-readable error message. This should be accurate but terse (typically one line). Always present.

D

Detail: an optional secondary error message carrying more detail about the problem. Might run to multiple lines.

H

Hint: an optional suggestion what to do about the problem. This is intended to differ from Detail in that it offers advice (potentially inappropriate) rather than hard facts. Might run to multiple lines.

P

Position: the field value is a decimal ASCII integer, indicating an error cursor position as an index into the original query string. The first character has index 1, and positions are measured in characters not bytes.
**Internal position:** this is defined the same as the `P` field, but it is used when the cursor position refers to an internally generated command rather than the one submitted by the client. The `q` field will always appear when this field appears.

**Internal query:** the text of a failed internally-generated command. This could be, for example, a SQL query issued by a PL/pgSQL function.

**Where:** an indication of the context in which the error occurred. Presently this includes a call stack traceback of active procedural language functions and internally-generated queries. The trace is one entry per line, most recent first.

**Schema name:** if the error was associated with a specific database object, the name of the schema containing that object, if any.

**Table name:** if the error was associated with a specific table, the name of the table. (Refer to the schema name field for the name of the table's schema.)

**Column name:** if the error was associated with a specific table column, the name of the column. (Refer to the schema and table name fields to identify the table.)

**Data type name:** if the error was associated with a specific data type, the name of the data type. (Refer to the schema name field for the name of the data type's schema.)

**Constraint name:** if the error was associated with a specific constraint, the name of the constraint. Refer to fields listed above for the associated table or domain. (For this purpose, indexes are treated as constraints, even if they weren't created with constraint syntax.)

**File:** the file name of the source-code location where the error was reported.

**Line:** the line number of the source-code location where the error was reported.

**Routine:** the name of the source-code routine reporting the error.

---

**Note**

The fields for schema name, table name, column name, data type name, and constraint name are supplied only for a limited number of error types; see Appendix A. Frontends should not assume that the presence of any of these fields guarantees the presence of another field. Core error sources observe the interrelationships noted above, but user-
defined functions may use these fields in other ways. In the same vein, clients should not assume that these fields denote contemporary objects in the current database.

The client is responsible for formatting displayed information to meet its needs; in particular it should break long lines as needed. Newline characters appearing in the error message fields should be treated as paragraph breaks, not line breaks.

### 53.9. Logical Replication Message Formats

This section describes the detailed format of each logical replication message. These messages are returned either by the replication slot SQL interface or are sent by a walsender. In case of a walsender they are encapsulated inside the replication protocol WAL messages as described in Section 53.4 and generally obey same message flow as physical replication.

#### Begin

- **Byte1('B')**
  - Identifies the message as a begin message.
- **Int64**
  - The final LSN of the transaction.
- **Int64**
  - Commit timestamp of the transaction. The value is in number of microseconds since PostgreSQL epoch (2000-01-01).
- **Int32**
  - Xid of the transaction.

#### Commit

- **Byte1('C')**
  - Identifies the message as a commit message.
- **Int8**
  - Flags; currently unused (must be 0).
- **Int64**
  - The LSN of the commit.
- **Int64**
  - The end LSN of the transaction.
- **Int64**
  - Commit timestamp of the transaction. The value is in number of microseconds since PostgreSQL epoch (2000-01-01).

#### Origin

- **Byte1('O')**
  - Identifies the message as an origin message.
Int64
   The LSN of the commit on the origin server.

String
   Name of the origin.

Note that there can be multiple Origin messages inside a single transaction.

Relation
   Byte1('R')
      Identifies the message as a relation message.

Int32
   ID of the relation.

String
   Namespace (empty string for pg_catalog).

String
   Relation name.

Int8
   Replica identity setting for the relation (same as relreplident in pg_class).

Int16
   Number of columns.

Next, the following message part appears for each column:

Int8
   Flags for the column. Currently can be either 0 for no flags or 1 which marks the column as part of the key.

String
   Name of the column.

Int32
   ID of the column’s data type.

Int32
   Type modifier of the column (atttypmod).

Type
   Byte1('Y')
      Identifies the message as a type message.

Int32
   ID of the data type.
String
  Namespace (empty string for \texttt{pg\_catalog}).

String
  Name of the data type.

Insert
  
  Byte1('I')
    Identifies the message as an insert message.

Int32
  ID of the relation corresponding to the ID in the relation message.

Byte1('N')
  Identifies the following TupleData message as a new tuple.

TupleData
  TupleData message part representing the contents of new tuple.

Update
  
  Byte1('U')
    Identifies the message as an update message.

Int32
  ID of the relation corresponding to the ID in the relation message.

Byte1('K')
  Identifies the following TupleData submessage as a key. This field is optional and is only present if the update changed data in any of the column(s) that are part of the REPLICA IDENTITY index.

Byte1('O')
  Identifies the following TupleData submessage as an old tuple. This field is optional and is only present if table in which the update happened has REPLICA IDENTITY set to FULL.

TupleData
  TupleData message part representing the contents of the old tuple or primary key. Only present if the previous 'O' or 'K' part is present.

Byte1('N')
  Identifies the following TupleData message as a new tuple.

TupleData
  TupleData message part representing the contents of a new tuple.

The Update message may contain either a 'K' message part or an 'O' message part or neither of them, but never both of them.
Delete

Byte1('D')

Identifies the message as a delete message.

Int32

ID of the relation corresponding to the ID in the relation message.

Byte1('K')

Identifies the following TupleData submessage as a key. This field is present if the table in which the delete has happened uses an index as REPLICA IDENTITY.

Byte1('O')

Identifies the following TupleData message as a old tuple. This field is present if the table in which the delete has happened has REPLICA IDENTITY set to FULL.

TupleData

TupleData message part representing the contents of the old tuple or primary key, depending on the previous field.

The Delete message may contain either a ‘K’ message part or an ‘O’ message part, but never both of them.

Truncate

Byte1('T')

Identifies the message as a truncate message.

Int32

Number of relations

Int8

Option bits for TRUNCATE: 1 for CASCADE, 2 for RESTART IDENTITY

Int32

ID of the relation corresponding to the ID in the relation message. This field is repeated for each relation.

The following message parts are shared by the above messages.

TupleData

Int16

Number of columns.

Next, one of the following submessages appears for each column:

Byte1('n')

Identifies the data as NULL value.

Or
Byte1('u')

Identifies unchanged TOASTed value (the actual value is not sent).

Or

Byte1('t')

Identifies the data as text formatted value.

Int32

Length of the column value.

Byte\n
The value of the column, in text format. (A future release might support additional formats.)

\n is the above length.

53.10. Summary of Changes since Protocol 2.0

This section provides a quick checklist of changes, for the benefit of developers trying to update existing client libraries to protocol 3.0.

The initial startup packet uses a flexible list-of-strings format instead of a fixed format. Notice that session default values for run-time parameters can now be specified directly in the startup packet. (Actually, you could do that before using the options field, but given the limited width of options and the lack of any way to quote whitespace in the values, it wasn't a very safe technique.)

All messages now have a length count immediately following the message type byte (except for startup packets, which have no type byte). Also note that PasswordMessage now has a type byte.

ErrorResponse and NoticeResponse ('E' and 'N') messages now contain multiple fields, from which the client code can assemble an error message of the desired level of verbosity. Note that individual fields will typically not end with a newline, whereas the single string sent in the older protocol always did.

The ReadyForQuery ('Z') message includes a transaction status indicator.

The distinction between BinaryRow and DataRow message types is gone; the single DataRow message type serves for returning data in all formats. Note that the layout of DataRow has changed to make it easier to parse. Also, the representation of binary values has changed: it is no longer directly tied to the server's internal representation.

There is a new “extended query” sub-protocol, which adds the frontend message types Parse, Bind, Execute, Describe, Close, Flush, and Sync, and the backend message types ParseComplete, BindComplete, PortalSuspended, ParameterDescription, NoData, and CloseComplete. Existing clients do not have to concern themselves with this sub-protocol, but making use of it might allow improvements in performance or functionality.

COPY data is now encapsulated into CopyData and CopyDone messages. There is a well-defined way to recover from errors during COPY. The special “\.” last line is not needed anymore, and is not sent during COPY OUT. (It is still recognized as a terminator during COPY IN, but its use is deprecated and will eventually be removed.) Binary COPY is supported. The CopyInResponse and CopyOutResponse messages include fields indicating the number of columns and the format of each column.

The layout of FunctionCall and FunctionCallResponse messages has changed. FunctionCall can now support passing NULL arguments to functions. It also can handle passing parameters and retrieving results in either text or binary format. There is no longer any reason to consider FunctionCall a potential security hole, since it does not offer direct access to internal server data representations.
The backend sends ParameterStatus (‘S’) messages during connection startup for all parameters it considers interesting to the client library. Subsequently, a ParameterStatus message is sent whenever the active value changes for any of these parameters.

The RowDescription (‘T’) message carries new table OID and column number fields for each column of the described row. It also shows the format code for each column.

The CursorResponse (‘P’) message is no longer generated by the backend.

The NotificationResponse (‘A’) message has an additional string field, which can carry a “payload” string passed from the NOTIFY event sender.

The EmptyQueryResponse (‘I’) message used to include an empty string parameter; this has been removed.
Chapter 54. PostgreSQL Coding Conventions

54.1. Formatting

Source code formatting uses 4 column tab spacing, with tabs preserved (i.e., tabs are not expanded to spaces). Each logical indentation level is one additional tab stop.

Layout rules (brace positioning, etc) follow BSD conventions. In particular, curly braces for the controlled blocks of if, while, switch, etc go on their own lines.

Limit line lengths so that the code is readable in an 80-column window. (This doesn't mean that you must never go past 80 columns. For instance, breaking a long error message string in arbitrary places just to keep the code within 80 columns is probably not a net gain in readability.)

Do not use C++ style comments ( // comments). Strict ANSI C compilers do not accept them. For the same reason, do not use C++ extensions such as declaring new variables mid-block.

The preferred style for multi-line comment blocks is

/*
 * comment text begins here
 * and continues here
 */

Note that comment blocks that begin in column 1 will be preserved as-is by pgindent, but it will reflow indented comment blocks as though they were plain text. If you want to preserve the line breaks in an indented block, add dashes like this:

/*----------
 * comment text begins here
 * and continues here
 ----------*/

While submitted patches do not absolutely have to follow these formatting rules, it's a good idea to do so. Your code will get run through pgindent before the next release, so there's no point in making it look nice under some other set of formatting conventions. A good rule of thumb for patches is “make the new code look like the existing code around it”.

The src/tools directory contains sample settings files that can be used with the emacs, xemacs or vim editors to help ensure that they format code according to these conventions.

The text browsing tools more and less can be invoked as:

```
more -x4
to make them show tabs appropriately.
```

54.2. Reporting Errors Within the Server
Error, warning, and log messages generated within the server code should be created using `ereport`, or its older cousin `elog`. The use of this function is complex enough to require some explanation.

There are two required elements for every message: a severity level (ranging from `DEBUG` to `PANIC`) and a primary message text. In addition there are optional elements, the most common of which is an error identifier code that follows the SQL spec's SQLSTATE conventions. `ereport` itself is just a shell function, that exists mainly for the syntactic convenience of making message generation look like a function call in the C source code. The only parameter accepted directly by `ereport` is the severity level. The primary message text and any optional message elements are generated by calling auxiliary functions, such as `errmsg`, within the `ereport` call.

A typical call to `ereport` might look like this:

```c
ereport(ERROR,
    (errcode(ERRCODE_DIVISION_BY_ZERO),
     errmsg("division by zero")));
```

This specifies error severity level `ERROR` (a run-of-the-mill error). The `errcode` call specifies the SQLSTATE error code using a macro defined in `src/include/utils/errcodes.h`. The `errmsg` call provides the primary message text. Notice the extra set of parentheses surrounding the auxiliary function calls — these are annoying but syntactically necessary.

Here is a more complex example:

```c
ereport(ERROR,
    (errcode(ERRCODE_AMBIGUOUS_FUNCTION),
     errmsg("function %s is not unique",
         func_signature_string(funcname, nargs, NIL, actual_arg_types)),
     errhint("Unable to choose a best candidate function. 
     "
     "You might need to add explicit typecasts.") ));
```

This illustrates the use of format codes to embed run-time values into a message text. Also, an optional “hint” message is provided.

If the severity level is `ERROR` or higher, `ereport` aborts the execution of the user-defined function and does not return to the caller. If the severity level is lower than `ERROR`, `ereport` returns normally.

The available auxiliary routines for `ereport` are:

- `errcode(sqlerrcode)` specifies the SQLSTATE error identifier code for the condition. If this routine is not called, the error identifier defaults to `ERRCODE_INTERNAL_ERROR` when the error severity level is `ERROR` or higher, `ERRCODE_WARNING` when the error level is `WARNING`, otherwise (for `NOTICE` and below) `ERRCODE_SUCCESSFUL_COMPLETION`. While these defaults are often convenient, always think whether they are appropriate before omitting the `errcode()` call.

- `errmsg(const char *msg, ...)` specifies the primary error message text, and possibly run-time values to insert into it. Insertions are specified by `sprintf`-style format codes. In addition to the standard format codes accepted by `sprintf`, the format code `%m` can be used to insert the error message returned by `strerror` for the current value of `errno`.¹ `%m` does not require any corresponding entry in the parameter list for `errmsg`. Note that the message string will be run through `gettext` for possible localization before format codes are processed.

- `errmsg_internal(const char *msg, ...)` is the same as `errmsg`, except that the message string will not be translated nor included in the internationalization message dictionary.

¹ That is, the value that was current when the `ereport` call was reached; changes of `errno` within the auxiliary reporting routines will not affect it. That would not be true if you were to write `strerror(errno)` explicitly in `errmsg`'s parameter list; accordingly, do not do so.
This should be used for “cannot happen” cases that are probably not worth expending translation effort on.

- `errmsg_plural` is like `errmsg`, but with support for various plural forms of the message. `fmt_singular` is the English singular format, `fmt_plural` is the English plural format, `n` is the integer value that determines which plural form is needed, and the remaining arguments are formatted according to the selected format string. For more information see Section 55.2.2.

- `errdetail` supplies an optional “detail” message; this is to be used when there is additional information that seems inappropriate to put in the primary message. The message string is processed in just the same way as for `errmsg`.

- `errdetail_internal` is the same as `errdetail`, except that the message string will not be translated nor included in the internationalization message dictionary. This should be used for detail messages that are not worth expending translation effort on, for instance because they are too technical to be useful to most users.

- `errdetail_plural` is like `errdetail`, but with support for various plural forms of the message. For more information see Section 55.2.2.

- `errdetail_log` is the same as `errdetail` except that this string goes only to the server log, never to the client. If both `errdetail` (or one of its equivalents above) and `errdetail_log` are used then one string goes to the client and the other to the log. This is useful for error details that are too security-sensitive or too bulky to include in the report sent to the client.

- `errdetail_log_plural` is like `errdetail_log`, but with support for various plural forms of the message. For more information see Section 55.2.2.

- `errhint` supplies an optional “hint” message; this is to be used when offering suggestions about how to fix the problem, as opposed to factual details about what went wrong. The message string is processed in just the same way as for `errmsg`.

- `errcontext` is not normally called directly from an `ereport` message site; rather it is used in `error_context_stack` callback functions to provide information about the context in which an error occurred, such as the current location in a PL function. The message string is processed in just the same way as for `errmsg`. Unlike the other auxiliary functions, this can be called more than once per `ereport` call; the successive strings thus supplied are concatenated with separating newlines.

- `errposition` specifies the textual location of an error within a query string. Currently it is only useful for errors detected in the lexical and syntactic analysis phases of query processing.

- `errtable` specifies a relation whose name and schema name should be included as auxiliary fields in the error report.

- `errtablecol` specifies a column whose name, table name, and schema name should be included as auxiliary fields in the error report.

- `errtableconstraint` specifies a table constraint whose name, table name, and schema name should be included as auxiliary fields in the error report. Indexes should be considered to be constraints for this purpose, whether or not they have an associated `pg_constraint` entry. Be careful to pass the underlying heap relation, not the index itself, as `rel`.

- `errdatatype` specifies a data type whose name and schema name should be included as auxiliary fields in the error report.
• `errdomainconstraint(Oid datatypeOid, const char *conname)` specifies a domain constraint whose name, domain name, and schema name should be included as auxiliary fields in the error report.

• `errcode_for_file_access()` is a convenience function that selects an appropriate SQLSTATE error identifier for a failure in a file-access-related system call. It uses the saved `errno` to determine which error code to generate. Usually this should be used in combination with `%m` in the primary error message text.

• `errcode_for_socket_access()` is a convenience function that selects an appropriate SQLSTATE error identifier for a failure in a socket-related system call.

• `errhidestmt(bool hide_stmt)` can be called to specify suppression of the `STATEMENT:` portion of a message in the postmaster log. Generally this is appropriate if the message text includes the current statement already.

• `errhidecontext(bool hide_ctx)` can be called to specify suppression of the `CONTEXT:` portion of a message in the postmaster log. This should only be used for verbose debugging messages where the repeated inclusion of context would bloat the log volume too much.

---

**Note**

At most one of the functions `errtable`, `errtablecol`, `errtableconstraint`, `errdatatype`, or `errdomainconstraint` should be used in an `ereport` call. These functions exist to allow applications to extract the name of a database object associated with the error condition without having to examine the potentially-localized error message text. These functions should be used in error reports for which it's likely that applications would wish to have automatic error handling. As of PostgreSQL 9.3, complete coverage exists only for errors in SQLSTATE class 23 (integrity constraint violation), but this is likely to be expanded in future.

There is an older function `elog` that is still heavily used. An `elog` call:

```c
elog(level, "format string", ...);
```

is exactly equivalent to:

```c
ereport(level, (errmsg_internal("format string", ...)));
```

Notice that the SQLSTATE error code is always defaulted, and the message string is not subject to translation. Therefore, `elog` should be used only for internal errors and low-level debug logging. Any message that is likely to be of interest to ordinary users should go through `ereport`. Nonetheless, there are enough internal “cannot happen” error checks in the system that `elog` is still widely used; it is preferred for those messages for its notational simplicity.

Advice about writing good error messages can be found in Section 54.3.

### 54.3. Error Message Style Guide

This style guide is offered in the hope of maintaining a consistent, user-friendly style throughout all the messages generated by PostgreSQL.

#### What Goes Where

The primary message should be short, factual, and avoid reference to implementation details such as specific function names. “Short” means “should fit on one line under normal conditions”. Use a detail
message if needed to keep the primary message short, or if you feel a need to mention implementation
details such as the particular system call that failed. Both primary and detail messages should be
factual. Use a hint message for suggestions about what to do to fix the problem, especially if the
suggestion might not always be applicable.

For example, instead of:

IpcMemoryCreate: shmget(key=%d, size=%u, 0%o) failed: %m
(plus a long addendum that is basically a hint)

write:

Primary: could not create shared memory segment: %m
Detail: Failed syscall was shmget(key=%d, size=%u, 0%o).
Hint: the addendum

Rationale: keeping the primary message short helps keep it to the point, and lets clients lay out screen
space on the assumption that one line is enough for error messages. Detail and hint messages can be
relegated to a verbose mode, or perhaps a pop-up error-details window. Also, details and hints would
normally be suppressed from the server log to save space. Reference to implementation details is best
avoided since users aren't expected to know the details.

Formatting

Don't put any specific assumptions about formatting into the message texts. Expect clients and the
server log to wrap lines to fit their own needs. In long messages, newline characters (\n) can be used
to indicate suggested paragraph breaks. Don't end a message with a newline. Don't use tabs or other
formatting characters. (In error context displays, newlines are automatically added to separate levels
of context such as function calls.)

Rationale: Messages are not necessarily displayed on terminal-type displays. In GUI displays or
browsers these formatting instructions are at best ignored.

Quotation Marks

English text should use double quotes when quoting is appropriate. Text in other languages should
consistently use one kind of quotes that is consistent with publishing customs and computer output
of other programs.

Rationale: The choice of double quotes over single quotes is somewhat arbitrary, but tends to be the
preferred use. Some have suggested choosing the kind of quotes depending on the type of object
according to SQL conventions (namely, strings single quoted, identifiers double quoted). But this is
a language-internal technical issue that many users aren't even familiar with, it won't scale to other
kinds of quoted terms, it doesn't translate to other languages, and it's pretty pointless, too.

Use of Quotes

Use quotes always to delimit file names, user-supplied identifiers, and other variables that might con-
tain words. Do not use them to mark up variables that will not contain words (for example, operator
names).

There are functions in the backend that will double-quote their own output at need (for example,
format_type_be()). Do not put additional quotes around the output of such functions.

Rationale: Objects can have names that create ambiguity when embedded in a message. Be consistent
about denoting where a plugged-in name starts and ends. But don't clutter messages with unnecessary
or duplicate quote marks.
Grammar and Punctuation

The rules are different for primary error messages and for detail/hint messages:

Primary error messages: Do not capitalize the first letter. Do not end a message with a period. Do not even think about ending a message with an exclamation point.

Detail and hint messages: Use complete sentences, and end each with a period. Capitalize the first word of sentences. Put two spaces after the period if another sentence follows (for English text; might be inappropriate in other languages).

Error context strings: Do not capitalize the first letter and do not end the string with a period. Context strings should normally not be complete sentences.

Rationale: Avoiding punctuation makes it easier for client applications to embed the message into a variety of grammatical contexts. Often, primary messages are not grammatically complete sentences anyway. (And if they're long enough to be more than one sentence, they should be split into primary and detail parts.) However, detail and hint messages are longer and might need to include multiple sentences. For consistency, they should follow complete-sentence style even when there's only one sentence.

Upper Case vs. Lower Case

Use lower case for message wording, including the first letter of a primary error message. Use upper case for SQL commands and key words if they appear in the message.

Rationale: It's easier to make everything look more consistent this way, since some messages are complete sentences and some not.

Avoid Passive Voice

Use the active voice. Use complete sentences when there is an acting subject (“A could not do B”). Use telegram style without subject if the subject would be the program itself; do not use “I” for the program.

Rationale: The program is not human. Don’t pretend otherwise.

Present vs. Past Tense

Use past tense if an attempt to do something failed, but could perhaps succeed next time (perhaps after fixing some problem). Use present tense if the failure is certainly permanent.

There is a nontrivial semantic difference between sentences of the form:

\[
\text{could not open file "}\%s\text{": }\%m
\]

and:

\[
\text{cannot open file "}\%s\text{"
}\]

The first one means that the attempt to open the file failed. The message should give a reason, such as “disk full” or “file doesn't exist”. The past tense is appropriate because next time the disk might not be full anymore or the file in question might exist.

The second form indicates that the functionality of opening the named file does not exist at all in the program, or that it’s conceptually impossible. The present tense is appropriate because the condition will persist indefinitely.
Rationale: Granted, the average user will not be able to draw great conclusions merely from the tense of the message, but since the language provides us with a grammar we should use it correctly.

**Type of the Object**

When citing the name of an object, state what kind of object it is.

Rationale: Otherwise no one will know what “foo.bar.baz” refers to.

**Brackets**

Square brackets are only to be used (1) in command synopses to denote optional arguments, or (2) to denote an array subscript.

Rationale: Anything else does not correspond to widely-known customary usage and will confuse people.

**Assembling Error Messages**

When a message includes text that is generated elsewhere, embed it in this style:

```
could not open file %s: %m
```

Rationale: It would be difficult to account for all possible error codes to paste this into a single smooth sentence, so some sort of punctuation is needed. Putting the embedded text in parentheses has also been suggested, but it's unnatural if the embedded text is likely to be the most important part of the message, as is often the case.

**Reasons for Errors**

Messages should always state the reason why an error occurred. For example:

```
BAD:    could not open file %s
BETTER: could not open file %s (I/O failure)
```

If no reason is known you better fix the code.

**Function Names**

Don’t include the name of the reporting routine in the error text. We have other mechanisms for finding that out when needed, and for most users it’s not helpful information. If the error text doesn’t make as much sense without the function name, reword it.

```
BAD:    pg_atoi: error in "z": cannot parse "z"
BETTER: invalid input syntax for integer: "z"
```

Avoid mentioning called function names, either; instead say what the code was trying to do:

```
BAD:    open() failed: %m
BETTER: could not open file %s: %m
```

If it really seems necessary, mention the system call in the detail message. (In some cases, providing the actual values passed to the system call might be appropriate information for the detail message.)

Rationale: Users don't know what all those functions do.
Tricky Words to Avoid

Unable. “Unable” is nearly the passive voice. Better use “cannot” or “could not”, as appropriate.

Bad. Error messages like “bad result” are really hard to interpret intelligently. It’s better to write why the result is “bad”, e.g., “invalid format”.

Illegal. “Illegal” stands for a violation of the law, the rest is “invalid”. Better yet, say why it’s invalid.

Unknown. Try to avoid “unknown”. Consider “error: unknown response”. If you don’t know what the response is, how do you know it’s erroneous? “Unrecognized” is often a better choice. Also, be sure to include the value being complained of.

BAD: unknown node type
BETTER: unrecognized node type: 42

Find vs. Exists. If the program uses a nontrivial algorithm to locate a resource (e.g., a path search) and that algorithm fails, it is fair to say that the program couldn’t “find” the resource. If, on the other hand, the expected location of the resource is known but the program cannot access it there then say that the resource doesn’t “exist”. Using “find” in this case sounds weak and confuses the issue.

May vs. Can vs. Might. “May” suggests permission (e.g., “You may borrow my rake.”), and has little use in documentation or error messages. “Can” suggests ability (e.g., “I can lift that log.”), and “might” suggests possibility (e.g., “It might rain today.”). Using the proper word clarifies meaning and assists translation.

Contractions. Avoid contractions, like “can’t”; use “cannot” instead.

Proper Spelling

Spell out words in full. For instance, avoid:

• spec
• stats
• parens
• auth
• xact

Rationale: This will improve consistency.

Localization

Keep in mind that error message texts need to be translated into other languages. Follow the guidelines in Section 55.2.2 to avoid making life difficult for translators.

54.4. Miscellaneous Coding Conventions

C Standard

Code in PostgreSQL should only rely on language features available in the C89 standard. That means a conforming C89 compiler has to be able to compile postgres, at least aside from a few platform
dependent pieces. Features from later revision of the C standard or compiler specific features can be used, if a fallback is provided.

For example static inline and _StaticAssert() are currently used, even though they are from newer revisions of the C standard. If not available we respectively fall back to defining the functions without inline, and to using a C89 compatible replacement that performs the same checks, but emits rather cryptic messages.

**Function-Like Macros and Inline Functions**

Both, macros with arguments and static inline functions, may be used. The latter are preferable if there are multiple-evaluation hazards when written as a macro, as e.g. the case with

```c
#define Max(x, y)       ((x) > (y) ? (x) : (y))
```

or when the macro would be very long. In other cases it's only possible to use macros, or at least easier. For example because expressions of various types need to be passed to the macro.

When the definition of an inline function references symbols (i.e. variables, functions) that are only available as part of the backend, the function may not be visible when included from frontend code.

```c
#ifndef FRONTEND
static inline MemoryContext
MemoryContextSwitchTo(MemoryContext context)
{
    MemoryContext old = CurrentMemoryContext;
    CurrentMemoryContext = context;
    return old;
}
#endif   /* FRONTEND */
```

In this example CurrentMemoryContext, which is only available in the backend, is referenced and the function thus hidden with a #ifndef FRONTEND. This rule exists because some compilers emit references to symbols contained in inline functions even if the function is not used.

**Writing Signal Handlers**

To be suitable to run inside a signal handler code has to be written very carefully. The fundamental problem is that, unless blocked, a signal handler can interrupt code at any time. If code inside the signal handler uses the same state as code outside chaos may ensue. As an example consider what happens if a signal handler tries to acquire a lock that's already held in the interrupted code.

Barring special arrangements code in signal handlers may only call async-signal safe functions (as defined in POSIX) and access variables of type volatile sig_atomic_t. A few functions in postgres are also deemed signal safe, importantly SetLatch().

In most cases signal handlers should do nothing more than note that a signal has arrived, and wake up code running outside of the handler using a latch. An example of such a handler is the following:

```c
static void
handle_sighup(SIGNAL_ARGS)
{
    int         save_errno = errno;

    got_SIGHUP = true;
}
```
SetLatch(MyLatch);

errno = save_errno;
}

errno is saved and restored because SetLatch() might change it. If that were not done interrupted code that's currently inspecting errno might see the wrong value.

**Calling Function Pointers**

For clarity, it is preferred to explicitly dereference a function pointer when calling the pointed-to function if the pointer is a simple variable, for example:

```c
(*emit_log_hook) (edata);
```

(even though `emit_log_hook(edata)` would also work). When the function pointer is part of a structure, then the extra punctuation can and usually should be omitted, for example:

```c
paramInfo->paramFetch(paramInfo, paramId);
```
Chapter 55. Native Language Support

55.1. For the Translator

PostgreSQL programs (server and client) can issue their messages in your favorite language — if the messages have been translated. Creating and maintaining translated message sets needs the help of people who speak their own language well and want to contribute to the PostgreSQL effort. You do not have to be a programmer at all to do this. This section explains how to help.

55.1.1. Requirements

We won't judge your language skills — this section is about software tools. Theoretically, you only need a text editor. But this is only in the unlikely event that you do not want to try out your translated messages. When you configure your source tree, be sure to use the --enable-nls option. This will also check for the libintl library and the msgfmt program, which all end users will need anyway. To try out your work, follow the applicable portions of the installation instructions.

If you want to start a new translation effort or want to do a message catalog merge (described later), you will need the programs xgettext and msgmerge, respectively, in a GNU-compatible implementation. Later, we will try to arrange it so that if you use a packaged source distribution, you won't need xgettext. (If working from Git, you will still need it.) GNU Gettext 0.10.36 or later is currently recommended.

Your local gettext implementation should come with its own documentation. Some of that is probably duplicated in what follows, but for additional details you should look there.

55.1.2. Concepts

The pairs of original (English) messages and their (possibly) translated equivalents are kept in message catalogs, one for each program (although related programs can share a message catalog) and for each target language. There are two file formats for message catalogs: The first is the “PO” file (for Portable Object), which is a plain text file with special syntax that translators edit. The second is the “MO” file (for Machine Object), which is a binary file generated from the respective PO file and is used while the internationalized program is run. Translators do not deal with MO files; in fact hardly anyone does.

The extension of the message catalog file is no surprise either .po or .mo. The base name is either the name of the program it accompanies, or the language the file is for, depending on the situation. This is a bit confusing. Examples are psql.po (PO file for psql) or fr.mo (MO file in French).

The file format of the PO files is illustrated here:

```
# comment
msgid "original string"
msgstr "translated string"
msgid "more original"
msgstr "another translated"
"string can be broken up like this"
...
```

The msgid's are extracted from the program source. (They need not be, but this is the most common way.) The msgstr lines are initially empty and are filled in with useful strings by the translator. The strings can contain C-style escape characters and can be continued across lines as illustrated. (The next line must start at the beginning of the line.)
Native Language Support

The # character introduces a comment. If whitespace immediately follows the # character, then this is a comment maintained by the translator. There can also be automatic comments, which have a non-whitespace character immediately following the #. These are maintained by the various tools that operate on the PO files and are intended to aid the translator.

#. automatic comment
#: filename.c:1023
#, flags, flags

The #. style comments are extracted from the source file where the message is used. Possibly the programmer has inserted information for the translator, such as about expected alignment. The #: comment indicates the exact location(s) where the message is used in the source. The translator need not look at the program source, but can if there is doubt about the correct translation. The #, comments contain flags that describe the message in some way. There are currently two flags: fuzzy is set if the message has possibly been outdated because of changes in the program source. The translator can then verify this and possibly remove the fuzzy flag. Note that fuzzy messages are not made available to the end user. The other flag is c-format, which indicates that the message is a printf-style format template. This means that the translation should also be a format string with the same number and type of placeholders. There are tools that can verify this, which key off the c-format flag.

55.1.3. Creating and Maintaining Message Catalogs

OK, so how does one create a “blank” message catalog? First, go into the directory that contains the program whose messages you want to translate. If there is a file nls.mk, then this program has been prepared for translation.

If there are already some .po files, then someone has already done some translation work. The files are named language.po, where language is the ISO 639-1 two-letter language code (in lower case), e.g., fr.po for French. If there is really a need for more than one translation effort per language then the files can also be named language_region.po where region is the ISO 3166-1 two-letter country code (in upper case), e.g., pt_BR.po for Portuguese in Brazil. If you find the language you wanted you can just start working on that file.

If you need to start a new translation effort, then first run the command:

make init-po

This will create a file proname.pot (.pot to distinguish it from PO files that are “in production”. The T stands for “template”.) Copy this file to language.po and edit it. To make it known that the new language is available, also edit the file nls.mk and add the language (or language and country) code to the line that looks like:

AVAIL LANGUAGES := de fr

(Other languages can appear, of course.)

As the underlying program or library changes, messages might be changed or added by the programmers. In this case you do not need to start from scratch. Instead, run the command:

make update-po

which will create a new blank message catalog file (the pot file you started with) and will merge it with the existing PO files. If the merge algorithm is not sure about a particular message it marks it “fuzzy” as explained above. The new PO file is saved with a .po.new extension.

2 https://www.iso.org/iso-3166-country-codes.html

2149
55.1.4. Editing the PO Files

The PO files can be edited with a regular text editor. The translator should only change the area between the quotes after the msgstr directive, add comments, and alter the fuzzy flag. There is (unsurprisingly) a PO mode for Emacs, which I find quite useful.

The PO files need not be completely filled in. The software will automatically fall back to the original string if no translation (or an empty translation) is available. It is no problem to submit incomplete translations for inclusions in the source tree; that gives room for other people to pick up your work. However, you are encouraged to give priority to removing fuzzy entries after doing a merge. Remember that fuzzy entries will not be installed; they only serve as reference for what might be the right translation.

Here are some things to keep in mind while editing the translations:

- Make sure that if the original ends with a newline, the translation does, too. Similarly for tabs, etc.
- If the original is a printf format string, the translation also needs to be. The translation also needs to have the same format specifiers in the same order. Sometimes the natural rules of the language make this impossible or at least awkward. In that case you can modify the format specifiers like this:

  ```
  msgstr "Die Datei %2$s hat %1$u Zeichen."
  ```

  Then the first placeholder will actually use the second argument from the list. The digits needs to follow the % immediately, before any other format manipulators. (This feature really exists in the printf family of functions. You might not have heard of it before because there is little use for it outside of message internationalization.)

- If the original string contains a linguistic mistake, report that (or fix it yourself in the program source) and translate normally. The corrected string can be merged in when the program sources have been updated. If the original string contains a factual mistake, report that (or fix it yourself) and do not translate it. Instead, you can mark the string with a comment in the PO file.

- Maintain the style and tone of the original string. Specifically, messages that are not sentences (cannot open file %s) should probably not start with a capital letter (if your language distinguishes letter case) or end with a period (if your language uses punctuation marks). It might help to read Section 54.3.

- If you don’t know what a message means, or if it is ambiguous, ask on the developers' mailing list. Chances are that English speaking end users might also not understand it or find it ambiguous, so it's best to improve the message.

55.2. For the Programmer

55.2.1. Mechanics

This section describes how to implement native language support in a program or library that is part of the PostgreSQL distribution. Currently, it only applies to C programs.

Adding NLS Support to a Program

1. Insert this code into the start-up sequence of the program:

   ```c
   #ifdef ENABLE_NLS
   #include <locale.h>
   #endif
   ```
Native Language Support

... 

#ifdef ENABLE_NLS
setlocale(LC_ALL, "");
bindtextdomain("progname", LOCALEDIR);
textdomain("progname");
#endif

(The progname can actually be chosen freely.)

2. Wherever a message that is a candidate for translation is found, a call to gettext() needs to be inserted. E.g.: 

fprintf(stderr, "panic level %d\n", lvl);

would be changed to:

fprintf(stderr, gettext("panic level %d\n"), lvl);

(gettext is defined as a no-op if NLS support is not configured.)

This tends to add a lot of clutter. One common shortcut is to use:

#define _(x) gettext(x)

Another solution is feasible if the program does much of its communication through one or a few functions, such as ereport() in the backend. Then you make this function call gettext internally on all input strings.

3. Add a file nls.mk in the directory with the program sources. This file will be read as a makefile. The following variable assignments need to be made here:

CATALOG_NAME

The program name, as provided in the textdomain() call.

AVAIL_LANGUAGES

List of provided translations — initially empty.

GETTEXT_FILES

List of files that contain translatable strings, i.e., those marked with gettext or an alternative solution. Eventually, this will include nearly all source files of the program. If this list gets too long you can make the first “file” be a + and the second word be a file that contains one file name per line.

GETTEXT_TRIGGERS

The tools that generate message catalogs for the translators to work on need to know what function calls contain translatable strings. By default, only gettext() calls are known. If you used _ or other identifiers you need to list them here. If the translatable string is not the first argument, the item needs to be of the form func:2 (for the second argument). If you have a function that supports pluralized messages, the item should look like func:1,2 (identifying the singular and plural message arguments).

The build system will automatically take care of building and installing the message catalogs.
55.2.2. Message-writing Guidelines

Here are some guidelines for writing messages that are easily translatable.

- Do not construct sentences at run-time, like:

```c
printf("Files were %s.\n", flag ? "copied" : "removed");
```

The word order within the sentence might be different in other languages. Also, even if you remember to call `gettext()` on each fragment, the fragments might not translate well separately. It's better to duplicate a little code so that each message to be translated is a coherent whole. Only numbers, file names, and such-like run-time variables should be inserted at run time into a message text.

- For similar reasons, this won't work:

```c
printf("copied %d file%s", n, n!=1 ? "s" : ");
```

because it assumes how the plural is formed. If you figured you could solve it like this:

```c
if (n==1)
    printf("copied 1 file");
else
    printf("copied %d files", n):
```

then be disappointed. Some languages have more than two forms, with some peculiar rules. It's often best to design the message to avoid the issue altogether, for instance like this:

```c
printf("number of copied files: %d", n);
```

If you really want to construct a properly pluralized message, there is support for this, but it's a bit awkward. When generating a primary or detail error message in `ereport()`, you can write something like this:

```c
errmsg_plural("copied %d file",
    "copied %d files",
    n,
    n)
```

The first argument is the format string appropriate for English singular form, the second is the format string appropriate for English plural form, and the third is the integer control value that determines which plural form to use. Subsequent arguments are formatted per the format string as usual. (Normally, the pluralization control value will also be one of the values to be formatted, so it has to be written twice.) In English it only matters whether `n` is 1 or not 1, but in other languages there can be many different plural forms. The translator sees the two English forms as a group and has the opportunity to supply multiple substitute strings, with the appropriate one being selected based on the run-time value of `n`.

If you need to pluralize a message that isn't going directly to an `errmsg` or `errdetail` report, you have to use the underlying function `ngettext`. See the gettext documentation.

- If you want to communicate something to the translator, such as about how a message is intended to line up with other output, precede the occurrence of the string with a comment that starts with `translator`, e.g.
/* translator: This message is not what it seems to be. */

These comments are copied to the message catalog files so that the translators can see them.
Chapter 56. Writing A Procedural Language Handler

All calls to functions that are written in a language other than the current “version 1” interface for compiled languages (this includes functions in user-defined procedural languages and functions written in SQL) go through a call handler function for the specific language. It is the responsibility of the call handler to execute the function in a meaningful way, such as by interpreting the supplied source text. This chapter outlines how a new procedural language’s call handler can be written.

The call handler for a procedural language is a “normal” function that must be written in a compiled language such as C, using the version-1 interface, and registered with PostgreSQL as taking no arguments and returning the type language_handler. This special pseudo-type identifies the function as a call handler and prevents it from being called directly in SQL commands. For more details on C language calling conventions and dynamic loading, see Section 38.10.

The call handler is called in the same way as any other function: It receives a pointer to a Function-CallInfoData struct containing argument values and information about the called function, and it is expected to return a Datum result (and possibly set the isnull field of the Function-CallInfoData structure, if it wishes to return an SQL null result). The difference between a call handler and an ordinary callee function is that the flinfo->fn_oid field of the FunctionCallInfoData structure will contain the OID of the actual function to be called, not of the call handler itself. The call handler must use this field to determine which function to execute. Also, the passed argument list has been set up according to the declaration of the target function, not of the call handler.

It’s up to the call handler to fetch the entry of the function from the pg_proc system catalog and to analyze the argument and return types of the called function. The AS clause from the CREATE FUNCTION command for the function will be found in the prosrc column of the pg_proc row. This is commonly source text in the procedural language, but in theory it could be something else, such as a path name to a file, or anything else that tells the call handler what to do in detail.

Often, the same function is called many times per SQL statement. A call handler can avoid repeated lookups of information about the called function by using the flinfo->fn_extra field. This will initially be NULL, but can be set by the call handler to point at information about the called function. On subsequent calls, if flinfo->fn_extra is already non-NULL then it can be used and the information lookup step skipped. The call handler must make sure that flinfo->fn_extra is made to point at memory that will live at least until the end of the current query, since an FmgrInfo data structure could be kept that long. One way to do this is to allocate the extra data in the memory context specified by flinfo->fn_mcxt: such data will normally have the same lifespan as the FmgrInfo itself. But the handler could also choose to use a longer-lived memory context so that it can cache function definition information across queries.

When a procedural-language function is invoked as a trigger, no arguments are passed in the usual way, but the FunctionCallInfoData's context field points at a TriggerData structure, rather than being NULL as it is in a plain function call. A language handler should provide mechanisms for procedural-language functions to get at the trigger information.

This is a template for a procedural-language handler written in C:

```c
#include "postgres.h"
#include "executor/spi.h"
#include "commands/trigger.h"
#include "fmgr.h"
#include "access/heapam.h"
#include "utils/syscache.h"
#include "catalog/pg_proc.h"
#include "catalog/pg_type.h"
```
PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(plsample_call_handler);

Datum
plsample_call_handler(PG_FUNCTION_ARGS)
{
    Datum          retval;

    if (CALLED_AS_TRIGGER(fcinfo))
    {
        /*
         * Called as a trigger function
         */
        TriggerData    *trigdata = (TriggerData *) fcinfo->context;
        retval = ...
    }
    else
    {
        /*
         * Called as a function
         */
        retval = ...
    }

    return retval;
}

Only a few thousand lines of code have to be added instead of the dots to complete the call handler.

After having compiled the handler function into a loadable module (see Section 38.10.5), the following
commands then register the sample procedural language:

CREATE FUNCTION plsample_call_handler() RETURNS language_handler
    AS 'filename'
    LANGUAGE C;
CREATE LANGUAGE plsample
    HANDLER plsample_call_handler;

Although providing a call handler is sufficient to create a minimal procedural language, there are two
other functions that can optionally be provided to make the language more convenient to use. These
are a validator and an inline handler. A validator can be provided to allow language-specific checking
to be done during CREATE FUNCTION. An inline handler can be provided to allow the language to
support anonymous code blocks executed via the DO command.

If a validator is provided by a procedural language, it must be declared as a function taking a single pa-
rameter of type oid. The validator's result is ignored, so it is customarily declared to return void. The
validator will be called at the end of a CREATE FUNCTION command that has created or updated a
function written in the procedural language. The passed-in OID is the OID of the function's pg_proc
row. The validator must fetch this row in the usual way, and do whatever checking is appropriate.
First, call CheckFunctionValidatorAccess() to diagnose explicit calls to the validator that
the user could not achieve through CREATE FUNCTION. Typical checks then include verifying that
the function's argument and result types are supported by the language, and that the function's body
is syntactically correct in the language. If the validator finds the function to be okay, it should just
return. If it finds an error, it should report that via the normal ereport() error reporting mechanism.
Throwing an error will force a transaction rollback and thus prevent the incorrect function definition from being committed.

Validator functions should typically honor the check_function_bodies parameter: if it is turned off then any expensive or context-sensitive checking should be skipped. If the language provides for code execution at compilation time, the validator must suppress checks that would induce such execution. In particular, this parameter is turned off by pg_dump so that it can load procedural language functions without worrying about side effects or dependencies of the function bodies on other database objects. (Because of this requirement, the call handler should avoid assuming that the validator has fully checked the function. The point of having a validator is not to let the call handler omit checks, but to notify the user immediately if there are obvious errors in a CREATE FUNCTION command.) While the choice of exactly what to check is mostly left to the discretion of the validator function, note that the core CREATE FUNCTION code only executes SET clauses attached to a function when check_function_bodies is on. Therefore, checks whose results might be affected by GUC parameters definitely should be skipped when check_function_bodies is off, to avoid false failures when reloading a dump.

If an inline handler is provided by a procedural language, it must be declared as a function taking a single parameter of type internal. The inline handler's result is ignored, so it is customarily declared to return void. The inline handler will be called when a DO statement is executed specifying the procedural language. The parameter actually passed is a pointer to an InlineCodeBlock struct, which contains information about the DO statement's parameters, in particular the text of the anonymous code block to be executed. The inline handler should execute this code and return.

It's recommended that you wrap all these function declarations, as well as the CREATE LANGUAGE command itself, into an extension so that a simple CREATE EXTENSION command is sufficient to install the language. See Section 38.16 for information about writing extensions.

The procedural languages included in the standard distribution are good references when trying to write your own language handler. Look into the src/pl subdirectory of the source tree. The CREATE LANGUAGE reference page also has some useful details.
Chapter 57. Writing A Foreign Data Wrapper

All operations on a foreign table are handled through its foreign data wrapper, which consists of a set of functions that the core server calls. The foreign data wrapper is responsible for fetching data from the remote data source and returning it to the PostgreSQL executor. If updating foreign tables is to be supported, the wrapper must handle that, too. This chapter outlines how to write a new foreign data wrapper.

The foreign data wrappers included in the standard distribution are good references when trying to write your own. Look into the contrib subdirectory of the source tree. The CREATE FOREIGN DATA WRAPPER reference page also has some useful details.

Note

The SQL standard specifies an interface for writing foreign data wrappers. However, PostgreSQL does not implement that API, because the effort to accommodate it into PostgreSQL would be large, and the standard API hasn't gained wide adoption anyway.

57.1. Foreign Data Wrapper Functions

The FDW author needs to implement a handler function, and optionally a validator function. Both functions must be written in a compiled language such as C, using the version-1 interface. For details on C language calling conventions and dynamic loading, see Section 38.10.

The handler function simply returns a struct of function pointers to callback functions that will be called by the planner, executor, and various maintenance commands. Most of the effort in writing an FDW is in implementing these callback functions. The handler function must be registered with PostgreSQL as taking no arguments and returning the special pseudo-type fdw_handler. The callback functions are plain C functions and are not visible or callable at the SQL level. The callback functions are described in Section 57.2.

The validator function is responsible for validating options given in CREATE and ALTER commands for its foreign data wrapper, as well as foreign servers, user mappings, and foreign tables using the wrapper. The validator function must be registered as taking two arguments, a text array containing the options to be validated, and an OID representing the type of object the options are associated with (in the form of the OID of the system catalog the object would be stored in, either ForeignDataWrapperRelationId, ForeignServerRelationId, UserMappingRelationId, or ForeignTableRelationId). If no validator function is supplied, options are not checked at object creation time or object alteration time.

57.2. Foreign Data Wrapper Callback Routines

The FDW handler function returns a palloc’d FdwRoutine struct containing pointers to the callback functions described below. The scan-related functions are required, the rest are optional.

The FdwRoutine struct type is declared in src/include/foreign/fdwapi.h, which see for additional details.

57.2.1. FDW Routines For Scanning Foreign Tables

2157
void GetForeignRelSize(PlannerInfo *root,
    RelOptInfo *baserel,
    Oid foreigntableid);

Obtain relation size estimates for a foreign table. This is called at the beginning of planning for a query that scans a foreign table. *root* is the planner's global information about the query; *baserel* is the planner's information about this table; and *foreigntableid* is the *pg_class* OID of the foreign table. (*foreigntableid* could be obtained from the planner data structures, but it's passed explicitly to save effort.)

This function should update *baserel-*->rows to be the expected number of rows returned by the table scan, after accounting for the filtering done by the restriction quals. The initial value of *baserel-*->rows is just a constant default estimate, which should be replaced if at all possible. The function may also choose to update *baserel-*->width if it can compute a better estimate of the average result row width.

See Section 57.4 for additional information.

void GetForeignPaths(PlannerInfo *root,
    RelOptInfo *baserel,
    Oid foreigntableid);

Create possible access paths for a scan on a foreign table. This is called during query planning. The parameters are the same as for GetForeignRelSize, which has already been called.

This function must generate at least one access path (*ForeignPath* node) for a scan on the foreign table and must call add_path to add each such path to *baserel-*->pathlist. It's recommended to use create_foreignscan_path to build the *ForeignPath* nodes. The function can generate multiple access paths, e.g., a path which has valid pathkeys to represent a pre-sorted result. Each access path must contain cost estimates, and can contain any FDW-private information that is needed to identify the specific scan method intended.

See Section 57.4 for additional information.

*ForeignScan* *
GetForeignPlan(PlannerInfo *root,
    RelOptInfo *baserel,
    Oid foreigntableid,
    ForeignPath *best_path,
    List *tlist,
    List *scan_clauses,
    Plan *outer_plan);

Create a *ForeignScan* plan node from the selected foreign access path. This is called at the end of query planning. The parameters are as for GetForeignRelSize, plus the selected *ForeignPath* (previously produced by GetForeignPaths, GetForeignJoinPaths, or GetForeignUpperPaths), the target list to be emitted by the plan node, the restriction clauses to be enforced by the plan node, and the outer subplan of the *ForeignScan*, which is used for rechecks performed by RecheckForeignScan. (If the path is for a join rather than a base relation, foreigntableid is InvalidOid.)

This function must create and return a *ForeignScan* plan node; it's recommended to use make_foreignscan to build the *ForeignScan* node.

See Section 57.4 for additional information.
void
BeginForeignScan(ForeignScanState *node, int eflags);

Begin executing a foreign scan. This is called during executor startup. It should perform any initial-
ization needed before the scan can start, but not start executing the actual scan (that should be done
upon the first call to IterateForeignScan). The ForeignScanState node has already been
created, but its fdw_state field is still NULL. Information about the table to scan is accessible
through the ForeignScanState node (in particular, from the underlying ForeignScan plan
node, which contains any FDW-private information provided by GetForeignPlan). eflags con-
tains flag bits describing the executor's operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not
perform any externally-visible actions; it should only do the minimum required to make the node state
valid for ExplainForeignScan and EndForeignScan.

TupleTableSlot *
IterateForeignScan(ForeignScanState *node);

Fetch one row from the foreign source, returning it in a tuple table slot (the node's ScanTupleSlot
should be used for this purpose). Return NULL if no more rows are available. The tuple table slot
infrastructure allows either a physical or virtual tuple to be returned; in most cases the latter choice is
preferable from a performance standpoint. Note that this is called in a short-lived memory context that
will be reset between invocations. Create a memory context in BeginForeignScan if you need
longer-lived storage, or use the es_query_cxt of the node's EState.

The rows returned must match the fdw_scan_tlist target list if one was supplied, otherwise they
must match the row type of the foreign table being scanned. If you choose to optimize away fetching
columns that are not needed, you should insert nulls in those column positions, or else generate a
fdw_scan_tlist list with those columns omitted.

Note that PostgreSQL's executor doesn't care whether the rows returned violate any constraints that
were defined on the foreign table — but the planner does care, and may optimize queries incorrectly
if there are rows visible in the foreign table that do not satisfy a declared constraint. If a constraint is
violated when the user has declared that the constraint should hold true, it may be appropriate to raise
an error (just as you would need to do in the case of a data type mismatch).

void
ReScanForeignScan(ForeignScanState *node);

Restart the scan from the beginning. Note that any parameters the scan depends on may have changed
value, so the new scan does not necessarily return exactly the same rows.

void
EndForeignScan(ForeignScanState *node);

End the scan and release resources. It is normally not important to release palloc'd memory, but for
example open files and connections to remote servers should be cleaned up.

57.2.2. FDW Routines For Scanning Foreign Joins

If an FDW supports performing foreign joins remotely (rather than by fetching both tables' data and
doing the join locally), it should provide this callback function:

void
GetForeignJoinPaths(PlannerInfo *root,
RelOptInfo *joinrel,
RelOptInfo *outerrel,
RelOptInfo *innerrel,
JoinType jointype,
JoinPathExtraData *extra);

Create possible access paths for a join of two (or more) foreign tables that all belong to the same foreign server. This optional function is called during query planning. As with GetForeignPaths, this function should generate ForeignPath path(s) for the supplied joinrel, and call add_path to add these paths to the set of paths considered for the join. But unlike GetForeignPaths, it is not necessary that this function succeed in creating at least one path, since paths involving local joining are always possible.

Note that this function will be invoked repeatedly for the same join relation, with different combinations of inner and outer relations; it is the responsibility of the FDW to minimize duplicated work.

If a ForeignPath path is chosen for the join, it will represent the entire join process; paths generated for the component tables and subsidiary joins will not be used. Subsequent processing of the join path proceeds much as it does for a path scanning a single foreign table. One difference is that the scanrelid of the resulting ForeignScan plan node should be set to zero, since there is no single relation that it represents; instead, the fs_relids field of the ForeignScan node represents the set of relations that were joined. (The latter field is set up automatically by the core planner code, and need not be filled by the FDW.) Another difference is that, because the column list for a remote join cannot be found from the system catalogs, the FDW must fill fdw_scan_tlist with an appropriate list of TargetEntry nodes, representing the set of columns it will supply at run time in the tuples it returns.

See Section 57.4 for additional information.

57.2.3. FDW Routines For Planning Post-Scan/Join Processing

If an FDW supports performing remote post-scan/join processing, such as remote aggregation, it should provide this callback function:

void
GetForeignUpperPaths(PlannerInfo *root,
                     UpperRelationKind stage,
                     RelOptInfo *input_rel,
                     RelOptInfo *output_rel,
                     void *extra);

Create possible access paths for upper relation processing, which is the planner’s term for all post-scan/join query processing, such as aggregation, window functions, sorting, and table updates. This optional function is called during query planning. Currently, it is called only if all base relation(s) involved in the query belong to the same FDW. This function should generate ForeignPath path(s) for any post-scan/join processing that the FDW knows how to perform remotely, and call add_path to add these paths to the indicated upper relation. As with GetForeignJoinPaths, it is not necessary that this function succeed in creating any paths, since paths involving local processing are always possible.

The stage parameter identifies which post-scan/join step is currently being considered. output_rel is the upper relation that should receive paths representing computation of this step, and input_rel is the relation representing the input to this step. The extra parameter provides additional details, currently, it is set only for UPPERREL_PARTIAL_GROUP_AGG or UPPERREL_GROUP_AGG, in which case it points to a GroupPathExtraData structure. (Note that ForeignPath paths added to output_rel would typically not have any direct dependency on paths of the input_rel, since their processing is expected to be done externally. However, examining paths previously generated for the previous processing step can be useful to avoid redundant planning work.)
See Section 57.4 for additional information.

57.2.4. FDW Routines For Updating Foreign Tables

If an FDW supports writable foreign tables, it should provide some or all of the following callback functions depending on the needs and capabilities of the FDW:

```c
void
AddForeignUpdateTargets(Query *parsetree,
    RangeTblEntry *target_rte,
    Relation target_relation);
```

UPDATE and DELETE operations are performed against rows previously fetched by the table-scanning functions. The FDW may need extra information, such as a row ID or the values of primary-key columns, to ensure that it can identify the exact row to update or delete. To support that, this function can add extra hidden, or “junk”, target columns to the list of columns that are to be retrieved from the foreign table during an UPDATE or DELETE.

To do that, add TargetEntry items to parsetree->targetList, containing expressions for the extra values to be fetched. Each such entry must be marked resjunk = true, and must have a distinct resname that will identify it at execution time. Avoid using names matching ctidN, wholerow, or wholerowN, as the core system can generate junk columns of these names. If the extra expressions are more complex than simple Vars, they must be run through eval_const_expressions before adding them to the targetlist.

Although this function is called during planning, the information provided is a bit different from that available to other planning routines. parsetree is the parse tree for the UPDATE or DELETE command, while target_rte and target_relation describe the target foreign table.

If the AddForeignUpdateTargets pointer is set to NULL, no extra target expressions are added. (This will make it impossible to implement DELETE operations, though UPDATE may still be feasible if the FDW relies on an unchanging primary key to identify rows.)

```c
List *
PlanForeignModify(PlannerInfo *root,
    ModifyTable *plan,
    Index resultRelation,
    int subplan_index);
```

Perform any additional planning actions needed for an insert, update, or delete on a foreign table. This function generates the FDW-private information that will be attached to the ModifyTable plan node that performs the update action. This private information must have the form of a List, and will be delivered to BeginForeignModify during the execution stage.

root is the planner's global information about the query. plan is the ModifyTable plan node, which is complete except for the fdwPrivLists field. resultRelation identifies the target foreign table by its range table index. subplan_index identifies which target of the ModifyTable plan node this is, counting from zero; use this if you want to index into plan->plans or other substructure of the plan node.

See Section 57.4 for additional information.

If the PlanForeignModify pointer is set to NULL, no additional plan-time actions are taken, and the fdw_private list delivered to BeginForeignModify will be NIL.
BeginForeignModify (ModifyTableState *mtstate,
    ResultRelInfo *rinfo,
    List *fdw_private,
    int subplan_index,
    int eflags);

Begin executing a foreign table modification operation. This routine is called during executor startup. It should perform any initialization needed prior to the actual table modifications. Subsequently, ExecForeignInsert, ExecForeignUpdate or ExecForeignDelete will be called for each tuple to be inserted, updated, or deleted.

mtstate is the overall state of the ModifyTable plan node being executed; global data about the plan and execution state is available via this structure. rinfo is the ResultRelInfo struct describing the target foreign table. (The ri_FdwState field of ResultRelInfo is available for the FDW to store any private state it needs for this operation.) fdw_private contains the private data generated by PlanForeignModify, if any. subplan_index identifies which target of the ModifyTable plan node this is. eflags contains flag bits describing the executor's operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not perform any externally-visible actions; it should only do the minimum required to make the node state valid for ExplainForeignModify and EndForeignModify.

If the BeginForeignModify pointer is set to NULL, no action is taken during executor startup.

TupleTableSlot *
ExecForeignInsert(EState *estate,
    ResultRelInfo *rinfo,
    TupleTableSlot *slot,
    TupleTableSlot *planSlot);

Insert one tuple into the foreign table. estate is global execution state for the query. rinfo is the ResultRelInfo struct describing the target foreign table. slot contains the tuple to be inserted; it will match the row-type definition of the foreign table. planSlot contains the tuple that was generated by the ModifyTable plan node's subplan; it differs from slot in possibly containing additional "junk" columns. (The planSlot is typically of little interest for INSERT cases, but is provided for completeness.)

The return value is either a slot containing the data that was actually inserted (this might differ from the data supplied, for example as a result of trigger actions), or NULL if no row was actually inserted (again, typically as a result of triggers). The passed-in slot can be re-used for this purpose.

The data in the returned slot is used only if the INSERT query has a RETURNING clause or the foreign table has an AFTER ROW trigger. Triggers require all columns, but the FDW could choose to optimize away returning some or all columns depending on the contents of the RETURNING clause. Regardless, some slot must be returned to indicate success, or the query's reported row count will be wrong.

If the ExecForeignInsert pointer is set to NULL, attempts to insert into the foreign table will fail with an error message.

TupleTableSlot *
ExecForeignUpdate(EState *estate,
    ResultRelInfo *rinfo,
    TupleTableSlot *slot,
    TupleTableSlot *planSlot);

Update one tuple in the foreign table. estate is global execution state for the query. rinfo is the ResultRelInfo struct describing the target foreign table. slot contains the new data for the tu-
Writing A Foreign Data Wrapper

...ple; it will match the row-type definition of the foreign table. planSlot contains the tuple that was
generated by the ModifyTable plan node's subplan; it differs from slot in possibly containing
additional "junk" columns. In particular, any junk columns that were requested by AddForeignUp-
dateTargets will be available from this slot.

The return value is either a slot containing the row as it was actually updated (this might differ from
the data supplied, for example as a result of trigger actions), or NULL if no row was actually updated
(again, typically as a result of triggers). The passed-in slot can be re-used for this purpose.

The data in the returned slot is used only if the UPDATE query has a RETURNING clause or the foreign
table has an AFTER ROW trigger. Triggers require all columns, but the FDW could choose to optimize
away returning some or all columns depending on the contents of the RETURNING clause. Regardless,
some slot must be returned to indicate success, or the query's reported row count will be wrong.

If the ExecForeignUpdate pointer is set to NULL, attempts to update the foreign table will fail
with an error message.

TupleTableSlot *
ExecForeignDelete(EState *estate,
    ResultRelInfo *rinfo,
    TupleTableSlot *slot,
    TupleTableSlot *planSlot);

Delete one tuple from the foreign table. estate is global execution state for the query. rinfo is the
ResultRelInfo struct describing the target foreign table. slot contains nothing useful upon call,
but can be used to hold the returned tuple. planSlot contains the tuple that was generated by the
ModifyTable plan node's subplan; in particular, it will carry any junk columns that were requested
by AddForeignUpdateTargets. The junk column(s) must be used to identify the tuple to be
deleted.

The return value is either a slot containing the row that was deleted, or NULL if no row was deleted
(typically as a result of triggers). The passed-in slot can be used to hold the tuple to be returned.

The data in the returned slot is used only if the DELETE query has a RETURNING clause or the foreign
table has an AFTER ROW trigger. Triggers require all columns, but the FDW could choose to optimize
away returning some or all columns depending on the contents of the RETURNING clause. Regardless,
some slot must be returned to indicate success, or the query's reported row count will be wrong.

If the ExecForeignDelete pointer is set to NULL, attempts to delete from the foreign table will fail
with an error message.

void
EndForeignModify(EState *estate,
    ResultRelInfo *rinfo);

End the table update and release resources. It is normally not important to release palloc'd memory,
but for example open files and connections to remote servers should be cleaned up.

If the EndForeignModify pointer is set to NULL, no action is taken during executor shutdown.

Tuples inserted into a partitioned table by INSERT or COPY FROM are routed to partitions. If an
FDW supports routable foreign-table partitions, it should also provide the following callback func-
tions. These functions are also called when COPY FROM is executed on a foreign table.

void
BeginForeignInsert(ModifyTableState *mtstate,
    ResultRelInfo *rinfo);
Begin executing an insert operation on a foreign table. This routine is called right before the first tuple is inserted into the foreign table in both cases when it is the partition chosen for tuple routing and the target specified in a COPY FROM command. It should perform any initialization needed prior to the actual insertion. Subsequently, ExecForeignInsert will be called for each tuple to be inserted into the foreign table.

mtstate is the overall state of the ModifyTable plan node being executed; global data about the plan and execution state is available via this structure. rinfo is the ResultRelInfo struct describing the target foreign table. (The ri_FdwState field of ResultRelInfo is available for the FDW to store any private state it needs for this operation.)

When this is called by a COPY FROM command, the plan-related global data in mtstate is not provided and the planSlot parameter of ExecForeignInsert subsequently called for each inserted tuple is NULL, whether the foreign table is the partition chosen for tuple routing or the target specified in the command.

If the BeginForeignInsert pointer is set to NULL, no action is taken for the initialization.

```void
EndForeignInsert(EState *estate,
                   ResultRelInfo *rinfo);
```

End the insert operation and release resources. It is normally not important to release palloc'd memory, but for example open files and connections to remote servers should be cleaned up.

If the EndForeignInsert pointer is set to NULL, no action is taken for the termination.

```int
IsForeignRelUpdatable(Relation rel);
```

Report which update operations the specified foreign table supports. The return value should be a bit mask of rule event numbers indicating which operations are supported by the foreign table, using the CmdType enumeration; that is, \((1 << \text{CMD\_UPDATE}) = 4\) for UPDATE, \((1 << \text{CMD\_INSERT}) = 8\) for INSERT, and \((1 << \text{CMD\_DELETE}) = 16\) for DELETE.

If the IsForeignRelUpdatable pointer is set to NULL, foreign tables are assumed to be insertable, updatable, or deletable if the FDW provides ExecForeignInsert, ExecForeignUpdate, or ExecForeignDelete respectively. This function is only needed if the FDW supports some tables that are updatable and some that are not. (Even then, it's permissible to throw an error in the execution routine instead of checking in this function. However, this function is used to determine updatability for display in the information_schema views.)

Some inserts, updates, and deletes to foreign tables can be optimized by implementing an alternative set of interfaces. The ordinary interfaces for inserts, updates, and deletes fetch rows from the remote server and then modify those rows one at a time. In some cases, this row-by-row approach is necessary, but it can be inefficient. If it is possible for the foreign server to determine which rows should be modified without actually retrieving them, and if there are no local triggers which would affect the operation, then it is possible to arrange things so that the entire operation is performed on the remote server. The interfaces described below make this possible.

```bool
PlanDirectModify(PlannerInfo *root,
                 ModifyTable *plan,
                 Index resultRelation,
                 int subplan_index);
```

Decide whether it is safe to execute a direct modification on the remote server. If so, return `true` after performing planning actions needed for that. Otherwise, return `false`. This optional function
is called during query planning. If this function succeeds, BeginDirectModify, IterateDirectModify and EndDirectModify will be called at the execution stage, instead. Otherwise, the table modification will be executed using the table-updating functions described above. The parameters are the same as for PlanForeignModify.

To execute the direct modification on the remote server, this function must rewrite the target subplan with a ForeignScan plan node that executes the direct modification on the remote server. The operation field of the ForeignScan must be set to the CmdType enumeration appropriately; that is, CMD_UPDATE for UPDATE, CMD_INSERT for INSERT, and CMD_DELETE for DELETE.

See Section 57.4 for additional information.

If the PlanDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.

```
void BeginDirectModify(ForeignScanState *node,
                      int eflags);
```

Prepare to execute a direct modification on the remote server. This is called during executor startup. It should perform any initialization needed prior to the direct modification (that should be done upon the first call to IterateDirectModify). The ForeignScanState node has already been created, but its fdw_state field is still NULL. Information about the table to modify is accessible through the ForeignScanState node (in particular, from the underlying ForeignScan plan node, which contains any FDW-private information provided by PlanDirectModify). eflags contains flag bits describing the executor's operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not perform any externally-visible actions; it should only do the minimum required to make the node state valid for ExplainDirectModify and EndDirectModify.

If the BeginDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.

```
TupleTableSlot *
IterateDirectModify(ForeignScanState *node);
```

When the INSERT, UPDATE or DELETE query doesn't have a RETURNING clause, just return NULL after a direct modification on the remote server. When the query has the clause, fetch one result containing the data needed for the RETURNING calculation, returning it in a tuple table slot (the node's ScanTupleSlot should be used for this purpose). The data that was actually inserted, updated or deleted must be stored in the es_result_relation_info->ri_projectReturning->pi_exprContext->ecxt_scantuple of the node's EState. Return NULL if no more rows are available. Note that this is called in a short-lived memory context that will be reset between invocations. Create a memory context in BeginDirectModify if you need longer-lived storage, or use the es_query_cxt of the node's EState.

The rows returned must match the fdw_scan_tlist target list if one was supplied, otherwise they must match the row type of the foreign table being updated. If you choose to optimize away fetching columns that are not needed for the RETURNING calculation, you should insert nulls in those column positions, or else generate a fdw_scan_tlist list with those columns omitted.

Whether the query has the clause or not, the query's reported row count must be incremented by the FDW itself. When the query doesn't have the clause, the FDW must also increment the row count for the ForeignScanState node in the EXPLAIN ANALYZE case.

If the IterateDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.
void
EndDirectModify(ForeignScanState *node);

Clean up following a direct modification on the remote server. It is normally not important to release
palloc'd memory, but for example open files and connections to the remote server should be cleaned up.

If the EndDirectModify pointer is set to NULL, no attempts to execute a direct modification on
the remote server are taken.

57.2.5. FDW Routines For Row Locking

If an FDW wishes to support late row locking (as described in Section 57.5), it must provide the
following callback functions:

RowMarkType
GetForeignRowMarkType(RangeTblEntry *rte,
    LockClauseStrength strength);

Report which row-marking option to use for a foreign table. rte is the RangeTblEntry node for
the table and strength describes the lock strength requested by the relevant FOR UPDATE/SHARE
clause, if any. The result must be a member of the RowMarkType enum type.

This function is called during query planning for each foreign table that appears in an UPDATE,
DELETE, or SELECT FOR UPDATE/SHARE query and is not the target of UPDATE or DELETE.

If the GetForeignRowMarkType pointer is set to NULL, the ROW_MARK_COPY option is always
used. (This implies that RefetchForeignRow will never be called, so it need not be provided
either.)

See Section 57.5 for more information.

HeapTuple
RefetchForeignRow(EState *estate,
    ExecRowMark *erm,
    Datum rowid,
    bool *updated);

Re-fetch one tuple from the foreign table, after locking it if required. estate is global execution state
for the query. erm is the ExecRowMark struct describing the target foreign table and the row lock
type (if any) to acquire. rowid identifies the tuple to be fetched. updated is an output parameter.

This function should return a palloc'ed copy of the fetched tuple, or NULL if the row lock couldn't be
obtained. The row lock type to acquire is defined by erm->markType, which is the value previously
returned by GetForeignRowMarkType. (ROW_MARK_REFERENCE means to just re-fetch the
tuple without acquiring any lock, and ROW_MARK_COPY will never be seen by this routine.)

In addition, *updated should be set to true if what was fetched was an updated version of the
tuple rather than the same version previously obtained. (If the FDW cannot be sure about this, always
returning true is recommended.)

Note that by default, failure to acquire a row lock should result in raising an error; a NULL return is
only appropriate if the SKIP LOCKED option is specified by erm->waitPolicy.

The rowid is the ctid value previously read for the row to be re-fetched. Although the rowid
value is passed as a Datum, it can currently only be a tid. The function API is chosen in hopes that
it may be possible to allow other data types for row IDs in future.
If the `RefetchForeignRow` pointer is set to NULL, attempts to re-fetch rows will fail with an error message.

See Section 57.5 for more information.

```c
bool
RecheckForeignScan(ForeignScanState *node,
                   TupleTableSlot *slot);
```

Recheck that a previously-returned tuple still matches the relevant scan and join qualifiers, and possibly provide a modified version of the tuple. For foreign data wrappers which do not perform join pushdown, it will typically be more convenient to set this to NULL and instead set `fdw_recheck_quals` appropriately. When outer joins are pushed down, however, it isn't sufficient to reapply the checks relevant to all the base tables to the result tuple, even if all needed attributes are present, because failure to match some qualifier might result in some attributes going to NULL, rather than in no tuple being returned. `RecheckForeignScan` can recheck qualifiers and return true if they are still satisfied and false otherwise, but it can also store a replacement tuple into the supplied slot.

To implement join pushdown, a foreign data wrapper will typically construct an alternative local join plan which is used only for rechecks: this will become the outer subplan of the `ForeignScan`. When a recheck is required, this subplan can be executed and the resulting tuple can be stored in the slot. This plan need not be efficient since no base table will return more than one row; for example, it may implement all joins as nested loops. The function `GetExistingLocalJoinPath` may be used to search existing paths for a suitable local join path, which can be used as the alternative local join plan. `GetExistingLocalJoinPath` searches for an unparameterized path in the path list of the specified join relation. (If it does not find such a path, it returns NULL, in which case a foreign data wrapper may build the local path by itself or may choose not to create access paths for that join.)

### 57.2.6. FDW Routines for EXPLAIN

```c
void
ExplainForeignScan(ForeignScanState *node,
                    ExplainState *es);
```

Print additional EXPLAIN output for a foreign table scan. This function can call `ExplainPropertyText` and related functions to add fields to the EXPLAIN output. The flag fields in `es` can be used to determine what to print, and the state of the `ForeignScanState` node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case.

If the `ExplainForeignScan` pointer is set to NULL, no additional information is printed during EXPLAIN.

```c
void
ExplainForeignModify(ModifyTableState *mtstate,
                      ResultRelInfo *rinfo,
                      List *fdw_private,
                      int subplan_index,
                      struct ExplainState *es);
```

Print additional EXPLAIN output for a foreign table update. This function can call `ExplainPropertyText` and related functions to add fields to the EXPLAIN output. The flag fields in `es` can be used to determine what to print, and the state of the `ModifyTableState` node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case. The first four arguments are the same as for `BeginForeignModify`.

If the `ExplainForeignModify` pointer is set to NULL, no additional information is printed during EXPLAIN.
void
ExplainDirectModify(ForeignScanState *node,
   ExplainState *es);

Print additional EXPLAIN output for a direct modification on the remote server. This function can call ExplainPropertyText and related functions to add fields to the EXPLAIN output. The flag fields in es can be used to determine what to print, and the state of the ForeignScanState node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case.

If the ExplainDirectModify pointer is set to NULL, no additional information is printed during EXPLAIN.

57.2.7. FDW Routines for ANALYZE

bool
AnalyzeForeignTable(Relation relation,
   AcquireSampleRowsFunc *func,
   BlockNumber *totalpages);

This function is called when ANALYZE is executed on a foreign table. If the FDW can collect statistics for this foreign table, it should return true, and provide a pointer to a function that will collect sample rows from the table in func, plus the estimated size of the table in pages in totalpages. Otherwise, return false.

If the FDW does not support collecting statistics for any tables, the AnalyzeForeignTable pointer can be set to NULL.

If provided, the sample collection function must have the signature

int
AcquireSampleRowsFunc(Relation relation,
   int elevel,
   HeapTuple *rows,
   int targrows,
   double *totalrows,
   double *totaldeadrows);

A random sample of up to targrows rows should be collected from the table and stored into the caller-provided rows array. The actual number of rows collected must be returned. In addition, store estimates of the total numbers of live and dead rows in the table into the output parameters totalrows and totaldeadrows. (Set totaldeadrows to zero if the FDW does not have any concept of dead rows.)

57.2.8. FDW Routines For IMPORT FOREIGN SCHEMA

List *
ImportForeignSchema(ImportForeignSchemaStmt *stmt, Oid serverOid);

Obtain a list of foreign table creation commands. This function is called when executing IMPORT FOREIGN SCHEMA, and is passed the parse tree for that statement, as well as the OID of the foreign server to use. It should return a list of C strings, each of which must contain a CREATE FOREIGN TABLE command. These strings will be parsed and executed by the core server.

Within the ImportForeignSchemaStmt struct, remote_schema is the name of the remote schema from which tables are to be imported. list_type identifies how to filter table names: FD-
Writing A Foreign Data Wrapper

W_IMPORT_SCHEMA_ALL means that all tables in the remote schema should be imported (in this case table_list is empty). FDW_IMPORT_SCHEMA_LIMIT_TO means to include only tables listed in table_list, and FDW_IMPORT_SCHEMA_EXCEPT means to exclude the tables listed in table_list. options is a list of options used for the import process. The meanings of the options are up to the FDW. For example, an FDW could use an option to define whether the NOT NULL attributes of columns should be imported. These options need not have anything to do with those supported by the FDW as database object options.

The FDW may ignore the local_schema field of the ImportForeignSchemaStmt, because the core server will automatically insert that name into the parsed CREATE FOREIGN TABLE commands.

The FDW does not have to concern itself with implementing the filtering specified by list_type and table_list, either, as the core server will automatically skip any returned commands for tables excluded according to those options. However, it's often useful to avoid the work of creating commands for excluded tables in the first place. The function IsImportableForeignTable() may be useful to test whether a given foreign-table name will pass the filter.

If the FDW does not support importing table definitions, the ImportForeignSchema pointer can be set to NULL.

57.2.9. FDW Routines for Parallel Execution

A ForeignScan node can, optionally, support parallel execution. A parallel ForeignScan will be executed in multiple processes and must return each row exactly once across all cooperating processes. To do this, processes can coordinate through fixed-size chunks of dynamic shared memory. This shared memory is not guaranteed to be mapped at the same address in every process, so it must not contain pointers. The following functions are all optional, but most are required if parallel execution is to be supported.

```c
bool IsForeignScanParallelSafe(PlannerInfo *root, RelOptInfo *rel, RangeTblEntry *rte);
```

Test whether a scan can be performed within a parallel worker. This function will only be called when the planner believes that a parallel plan might be possible, and should return true if it is safe for that scan to run within a parallel worker. This will generally not be the case if the remote data source has transaction semantics, unless the worker's connection to the data can somehow be made to share the same transaction context as the leader.

If this function is not defined, it is assumed that the scan must take place within the parallel leader. Note that returning true does not mean that the scan itself can be done in parallel, only that the scan can be performed within a parallel worker. Therefore, it can be useful to define this method even when parallel execution is not supported.

```c
Size EstimateDSMForeignScan(ForeignScanState *node, ParallelContext *pcxt);
```

Estimate the amount of dynamic shared memory that will be required for parallel operation. This may be higher than the amount that will actually be used, but it must not be lower. The return value is in bytes. This function is optional, and can be omitted if not needed; but if it is omitted, the next three functions must be omitted as well, because no shared memory will be allocated for the FDW’s use.

```c
void InitializeDSMForeignScan(ForeignScanState *node, ParallelContext *pcxt,
```
void *coordinate);

Initialize the dynamic shared memory that will be required for parallel operation. coordinate points to a shared memory area of size equal to the return value of EstimateDSMForeignScan. This function is optional, and can be omitted if not needed.

void ReInitializeDSMForeignScan(ForeignScanState *node, ParallelContext *pcxt,
                                 void *coordinate);

Re-initialize the dynamic shared memory required for parallel operation when the foreign-scan plan node is about to be re-scanned. This function is optional, and can be omitted if not needed. Recommended practice is that this function reset only shared state, while the ReScanForeignScan function resets only local state. Currently, this function will be called before ReScanForeignScan, but it's best not to rely on that ordering.

void InitializeWorkerForeignScan(ForeignScanState *node, shm_toc *toc,
                                 void *coordinate);

Initialize a parallel worker's local state based on the shared state set up by the leader during InitializeDSMForeignScan. This function is optional, and can be omitted if not needed.

void ShutdownForeignScan(ForeignScanState *node);

Release resources when it is anticipated the node will not be executed to completion. This is not called in all cases; sometimes, EndForeignScan may be called without this function having been called first. Since the DSM segment used by parallel query is destroyed just after this callback is invoked, foreign data wrappers that wish to take some action before the DSM segment goes away should implement this method.

57.2.10. FDW Routines For reparameterization of paths

List *
ReparameterizeForeignPathByChild(PlannerInfo *root, List *fdw_private,
                                 RelOptInfo *child_rel);

This function is called while converting a path parameterized by the top-most parent of the given child relation child_rel to be parameterized by the child relation. The function is used to reparameterize any paths or translate any expression nodes saved in the given fdw_private member of a ForeignPath. The callback may use reparameterize_path_by_child, adjust_appendrel_attrs or adjust_appendrel_attrs_multilevel as required.

57.3. Foreign Data Wrapper Helper Functions

Several helper functions are exported from the core server so that authors of foreign data wrappers can get easy access to attributes of FDW-related objects, such as FDW options. To use any of these functions, you need to include the header file foreign/foreign.h in your source file. That header also defines the struct types that are returned by these functions.
ForeignDataWrapper *
GetForeignDataWrapper(Oid fdwid);

This function returns a ForeignDataWrapper object for the foreign-data wrapper with the given OID. A ForeignDataWrapper object contains properties of the FDW (see foreign/foreign.h for details).

ForeignServer *
GetForeignServer(Oid serverid);

This function returns a ForeignServer object for the foreign server with the given OID. A ForeignServer object contains properties of the server (see foreign/foreign.h for details).

UserMapping *
GetUserMapping(Oid userid, Oid serverid);

This function returns a UserMapping object for the user mapping of the given role on the given server. (If there is no mapping for the specific user, it will return the mapping for PUBLIC, or throw error if there is none.) A UserMapping object contains properties of the user mapping (see foreign/foreign.h for details).

ForeignTable *
GetForeignTable(Oid relid);

This function returns a ForeignTable object for the foreign table with the given OID. A ForeignTable object contains properties of the foreign table (see foreign/foreign.h for details).

List *
GetForeignColumnOptions(Oid relid, AttrNumber attnum);

This function returns the per-column FDW options for the column with the given foreign table OID and attribute number, in the form of a list of DefElem. NIL is returned if the column has no options.

Some object types have name-based lookup functions in addition to the OID-based ones:

ForeignDataWrapper *
GetForeignDataWrapperByName(const char *name, bool missing_ok);

This function returns a ForeignDataWrapper object for the foreign-data wrapper with the given name. If the wrapper is not found, return NULL if missing_ok is true, otherwise raise an error.

ForeignServer *
GetForeignServerByName(const char *name, bool missing_ok);

This function returns a ForeignServer object for the foreign server with the given name. If the server is not found, return NULL if missing_ok is true, otherwise raise an error.

57.4. Foreign Data Wrapper Query Planning

The FDW callback functions GetForeignRelSize, GetForeignPaths, GetForeignPlan, PlanForeignModify, GetForeignJoinPaths, GetForeignUpperPaths, and PlanDirectModify must fit into the workings of the PostgreSQL planner. Here are some notes about what they must do.
The information in `root` and `baserel` can be used to reduce the amount of information that has to be fetched from the foreign table (and therefore reduce the cost). `baserel->baserestrictinfo` is particularly interesting, as it contains restriction quals (WHERE clauses) that should be used to filter the rows to be fetched. (The FDW itself is not required to enforce these quals, as the core executor can check them instead.) `baserel->reltarget->exprs` can be used to determine which columns need to be fetched; but note that it only lists columns that have to be emitted by the `ForeignScan` plan node, not columns that are used in qual evaluation but not output by the query.

Various private fields are available for the FDW planning functions to keep information in. Generally, whatever you store in FDW private fields should be `palloc`d, so that it will be reclaimed at the end of planning.

`baserel->fdw_private` is a `void` pointer that is available for FDW planning functions to store information relevant to the particular foreign table. The core planner does not touch it except to initialize it to NULL when the `RelOptInfo` node is created. It is useful for passing information forward from `GetForeignRelSize` to `GetForeignPaths` and/or `GetForeignPaths` to `GetForeignPlan`, thereby avoiding recalculation.

`GetForeignPaths` can identify the meaning of different access paths by storing private information in the `fdw_private` field of `ForeignPath` nodes. `fdw_private` is declared as a `List` pointer, but could actually contain anything since the core planner does not touch it. However, best practice is to use a representation that's dumpable by `nodeToString`, for use with debugging support available in the backend.

`GetForeignPlan` can examine the `fdw_private` field of the selected `ForeignPath` node, and can generate `fdw_exprs` and `fdw_private` lists to be placed in the `ForeignScan` plan node, where they will be available at execution time. Both of these lists must be represented in a form that `copyObject` knows how to copy. The `fdw_private` list has no other restrictions and is not interpreted by the core backend in any way. The `fdw_exprs` list, if not NIL, is expected to contain expression trees that are intended to be executed at run time. These trees will undergo post-processing by the planner to make them fully executable.

In `GetForeignPlan`, generally the passed-in target list can be copied into the plan node as-is. The passed `scan_clauses` list contains the same clauses as `baserel->baserestrictinfo`, but may be re-ordered for better execution efficiency. In simple cases the FDW can just strip `RestrictInfo` nodes from the `scan_clauses` list (using `extract_actual_clauses`) and put all the clauses into the plan node's qual list, which means that all the clauses will be checked by the executor at run time. More complex FDWs may be able to check some of the clauses internally, in which case those clauses can be removed from the plan node's qual list so that the executor doesn't waste time rechecking them.

As an example, the FDW might identify some restriction clauses of the form `foreign_variable = sub_expression`, which it determines can be executed on the remote server given the locally-evaluated value of the `sub_expression`. The actual identification of such a clause should happen during `GetForeignPaths`, since it would affect the cost estimate for the path. The path's `fdw_private` field would probably include a pointer to the identified clause's `RestrictInfo` node. Then `GetForeignPlan` would remove that clause from `scan_clauses`, but add the `sub_expression` to `fdw_exprs` to ensure that it gets massaged into executable form. It would probably also put control information into the plan node's `fdw_private` field to tell the execution functions what to do at run time. The query transmitted to the remote server would involve something like `WHERE foreign_variable = $1`, with the parameter value obtained at run time from evaluation of the `fdw_exprs` expression tree.

Any clauses removed from the plan node's qual list must instead be added to `fdw_recheck_quals` or rechecked by `RecheckForeignScan` in order to ensure correct behavior at the READ COMMITTED isolation level. When a concurrent update occurs for some other table involved in the query, the executor may need to verify that all of the original quals are still satisfied for the tuple, possibly against a different set of parameter values. Using `fdw_recheck_quals` is typically easier than implementing checks inside `RecheckForeignScan`, but this method will be insufficient when outer
joins have been pushed down, since the join tuples in that case might have some fields go to NULL without rejecting the tuple entirely.

Another ForeignScan field that can be filled by FDWs is fdw_scan_tlist, which describes the tuples returned by the FDW for this plan node. For simple foreign table scans this can be set to NIL, implying that the returned tuples have the row type declared for the foreign table. A non-NIL value must be a target list (list of TargetEntries) containing Vars and/or expressions representing the returned columns. This might be used, for example, to show that the FDW has omitted some columns that it noticed won’t be needed for the query. Also, if the FDW can compute expressions used by the query more cheaply than can be done locally, it could add those expressions to fdw_scan_tlist. Note that join plans (created from paths made by GetForeignJoinPaths) must always supply fdw_scan_tlist to describe the set of columns they will return.

The FDW should always construct at least one path that depends only on the table’s restriction clauses. In join queries, it might also choose to construct path(s) that depend on join clauses, for example foreign_variable = local_variable. Such clauses will not be found in baserel->baserestrictinfo but must be sought in the relation’s join lists. A path using such a clause is called a “parameterized path”. It must identify the other relations used in the selected join clause(s) with a suitable value of param_info; use get_baserel_parampathinfo to compute that value. In GetForeignPlan, the local_variable portion of the join clause would be added to fdw_exprs, and then at run time the case works the same as for an ordinary restriction clause.

If an FDW supports remote joins, GetForeignJoinPaths should produce ForeignPaths for potential remote joins in much the same way as GetForeignPaths works for base tables. Information about the intended join can be passed forward to GetForeignPlan in the same ways described above. However, baserestrictinfo is not relevant for join relations; instead, the relevant join clauses for a particular join are passed to GetForeignJoinPaths as a separate parameter (extra->restrictlist).

An FDW might additionally support direct execution of some plan actions that are above the level of scans and joins, such as grouping or aggregation. To offer such options, the FDW should generate paths and insert them into the appropriate upper relation. For example, a path representing remote aggregation should be inserted into the UPPERREL_GROUP_AGG relation, using add_path. This path will be compared on a cost basis with local aggregation performed by reading a simple scan path for the foreign relation (note that such a path must also be supplied, else there will be an error at plan time). If the remote-aggregation path wins, which it usually would, it will be converted into a plan in the usual way, by calling GetForeignPlan. The recommended place to generate such paths is in the GetForeignUpperPaths callback function, which is called for each upper relation (i.e., each post-scan/join processing step), if all the base relations of the query come from the same FDW.

PlanForeignModify and the other callbacks described in Section 57.2.4 are designed around the assumption that the foreign relation will be scanned in the usual way and then individual row updates will be driven by a local ModifyTable plan node. This approach is necessary for the general case where an update requires reading local tables as well as foreign tables. However, if the operation could be executed entirely by the foreign server, the FDW could generate a path representing that and insert it into the UPPERREL_FINAL upper relation, where it would compete against the ModifyTable approach. This approach could also be used to implement remote SELECT FOR UPDATE, rather than using the row locking callbacks described in Section 57.2.5. Keep in mind that a path inserted into UPPERREL_FINAL is responsible for implementing all behavior of the query.

When planning an UPDATE or DELETE, PlanForeignModify and PlanDirectModify can look up the RelOptInfo struct for the foreign table and make use of the baserel->fdw_private data previously created by the scan-planning functions. However, in INSERT the target table is not scanned so there is no RelOptInfo for it. The List returned by PlanForeignModify has the same restrictions as the fdw_private list of a ForeignScan plan node, that is it must contain only structures that copyObject knows how to copy.

INSERT with an ON CONFLICT clause does not support specifying the conflict target, as unique constraints or exclusion constraints on remote tables are not locally known. This in turn implies that ON CONFLICT DO UPDATE is not supported, since the specification is mandatory there.
57.5. Row Locking in Foreign Data Wrappers

If an FDW’s underlying storage mechanism has a concept of locking individual rows to prevent concurrent updates of those rows, it is usually worthwhile for the FDW to perform row-level locking with as close an approximation as practical to the semantics used in ordinary PostgreSQL tables. There are multiple considerations involved in this.

One key decision to be made is whether to perform *early locking* or *late locking*. In early locking, a row is locked when it is first retrieved from the underlying store, while in late locking, the row is locked only when it is known that it needs to be locked. (The difference arises because some rows may be discarded by locally-checked restriction or join conditions.) Early locking is much simpler and avoids extra round trips to a remote store, but it can cause locking of rows that need not have been locked, resulting in reduced concurrency or even unexpected deadlocks. Also, late locking is only possible if the row to be locked can be uniquely re-identified later. Preferably the row identifier should identify a specific version of the row, as PostgreSQL TIDs do.

By default, PostgreSQL ignores locking considerations when interfacing to FDWs, but an FDW can perform early locking without any explicit support from the core code. The API functions described in Section 57.2.5, which were added in PostgreSQL 9.5, allow an FDW to use late locking if it wishes.

An additional consideration is that in *READ COMMITTED* isolation mode, PostgreSQL may need to re-check restriction and join conditions against an updated version of some target tuple. Rechecking join conditions requires re-obtaining copies of the non-target rows that were previously joined to the target tuple. When working with standard PostgreSQL tables, this is done by including the TIDs of the non-target tables in the column list projected through the join, and then re-fetching non-target rows when required. This approach keeps the join data set compact, but it requires inexpensive re-fetch capability, as well as a TID that can uniquely identify the row version to be re-fetched. By default, therefore, the approach used with foreign tables is to include a copy of the entire row fetched from a foreign table in the column list projected through the join. This puts no special demands on the FDW but can result in reduced performance of merge and hash joins. An FDW that is capable of meeting the re-fetch requirements can choose to do it the first way.

For an *UPDATE* or *DELETE* on a foreign table, it is recommended that the *ForeignScan* operation on the target table perform early locking on the rows that it fetches, perhaps via the equivalent of *SELECT FOR UPDATE*. An FDW can detect whether a table is an *UPDATE*/*DELETE* target at plan time by comparing its relid to `root->parse->resultRelation`, or at execution time by using `ExecRelationIsTargetRelation()`. An alternative possibility is to perform late locking within the `ExecForeignUpdate` or `ExecForeignDelete` callback, but no special support is provided for this.

For foreign tables that are specified to be locked by a *SELECT FOR UPDATE/SHARE* command, the *ForeignScan* operation can again perform early locking by fetching tuples with the equivalent of *SELECT FOR UPDATE/SHARE*. To perform late locking instead, provide the callback functions defined in Section 57.2.5. In `GetForeignRowMarkType`, select rowmark option `ROW_MARK_EXCLUSIVE, ROW_MARK_NOKEYEXCLUSIVE, ROW_MARK_SHARE, or ROW_MARK_KEYSHARE` depending on the requested lock strength. (The core code will act the same regardless of which of these four options you choose.) Elsewhere, you can detect whether a foreign table was specified to be locked by this type of command by using `get_plan_rowmark` at plan time, or `ExecFindRowMark` at execution time; you must check not only whether a non-null rowmark struct is returned, but that its `strength` field is not `LCS_NONE`.

Lastly, for foreign tables that are used in an *UPDATE*, *DELETE* or *SELECT FOR UPDATE/SHARE* command but are not specified to be row-locked, you can override the default choice to copy entire rows by having `GetForeignRowMarkType` select option `ROW_MARK_REFERENCE` when it sees lock strength `LCS_NONE`. This will cause `RefetchForeignRow` to be called with that value for `markType`; it should then re-fetch the row without acquiring any new lock. (If you have a `GetForeignRowMarkType` function but don’t wish to re-fetch unlocked rows, select option `ROW_MARK_COPY` for `LCS_NONE`.)
See src/include/nodes/lockoptions.h, the comments for RowMarkType and PlanRowMark in src/include/nodes/plannodes.h, and the comments for ExecRowMark in src/include/nodes/execnodes.h for additional information.
Chapter 58. Writing A Table Sampling Method

PostgreSQL’s implementation of the TABLESAMPLE clause supports custom table sampling methods, in addition to the BERNOULLI and SYSTEM methods that are required by the SQL standard. The sampling method determines which rows of the table will be selected when the TABLESAMPLE clause is used.

At the SQL level, a table sampling method is represented by a single SQL function, typically implemented in C, having the signature

```
method_name(internal) RETURNS tsm_handler
```

The name of the function is the same method name appearing in the TABLESAMPLE clause. The internal argument is a dummy (always having value zero) that simply serves to prevent this function from being called directly from a SQL command. The result of the function must be a palloc’d struct of type TsmRoutine, which contains pointers to support functions for the sampling method. These support functions are plain C functions and are not visible or callable at the SQL level. The support functions are described in Section 58.1.

In addition to function pointers, the TsmRoutine struct must provide these additional fields:

- **List *parameterTypes**
  
  This is an OID list containing the data type OIDs of the parameter(s) that will be accepted by the TABLESAMPLE clause when this sampling method is used. For example, for the built-in methods, this list contains a single item with value FLOAT4OID, which represents the sampling percentage. Custom sampling methods can have more or different parameters.

- **bool repeatable_across_queries**
  
  If true, the sampling method can deliver identical samples across successive queries, if the same parameters and REPEATABLE seed value are supplied each time and the table contents have not changed. When this is false, the REPEATABLE clause is not accepted for use with the sampling method.

- **bool repeatable_across_scans**
  
  If true, the sampling method can deliver identical samples across successive scans in the same query (assuming unchanging parameters, seed value, and snapshot). When this is false, the planner will not select plans that would require scanning the sampled table more than once, since that might result in inconsistent query output.

The TsmRoutine struct type is declared in src/include/access/tsmapi.h, which see for additional details.

The table sampling methods included in the standard distribution are good references when trying to write your own. Look into the src/backend/access/tablesample subdirectory of the source tree for the built-in sampling methods, and into the contrib subdirectory for add-on methods.

58.1. Sampling Method Support Functions

The TSM handler function returns a palloc’d TsmRoutine struct containing pointers to the support functions described below. Most of the functions are required, but some are optional, and those pointers can be NULL.
void SampleScanGetSampleSize (PlannerInfo *root,
   RelOptInfo *baserel,
   List *paramexprs,
   BlockNumber *pages,
   double *tuples);

This function is called during planning. It must estimate the number of relation pages that will be
read during a sample scan, and the number of tuples that will be selected by the scan. (For example,
these might be determined by estimating the sampling fraction, and then multiplying the baserel->pages and baserel->tuples numbers by that, being sure to round the results to integral
values.) The paramexprs list holds the expression(s) that are parameters to the TABLESAMPLE
clause. It is recommended to use estimate_expression_value() to try to reduce these ex-
pressions to constants, if their values are needed for estimation purposes; but the function must provide
size estimates even if they cannot be reduced, and it should not fail even if the values appear invalid
(remember that they're only estimates of what the run-time values will be). The pages and tuples
parameters are outputs.

void InitSampleScan (SampleScanState *node,
   int eflags);

Initialize for execution of a SampleScan plan node. This is called during executor startup. It should
perform any initialization needed before processing can start. The SampleScanState node has al-
ready been created, but its tsm_state field is NULL. The InitSampleScan function can palloc
whatever internal state data is needed by the sampling method, and store a pointer to it in node->
tsm_state. Information about the table to scan is accessible through other fields of the Sam-
pleScanState node (but note that the node->ss.ss_currentScanDesc scan descriptor is
not set up yet). eflags contains flag bits describing the executor's operating mode for this plan node.

When (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, the scan will not actually be performed,
so this function should only do the minimum required to make the node state valid for EXPLAIN and
EndSampleScan.

This function can be omitted (set the pointer to NULL), in which case BeginSampleScan must
perform all initialization needed by the sampling method.

void BeginSampleScan (SampleScanState *node,
   Datum *params,
   int nparams,
   uint32 seed);

Begin execution of a sampling scan. This is called just before the first attempt to fetch a tuple, and
may be called again if the scan needs to be restarted. Information about the table to scan is accessi-
ble through fields of the SampleScanState node (but note that the node->ss.ss_currentScanDesc scan descriptor is not set up yet). The params array, of length nparams, contains the
values of the parameters supplied in the TABLESAMPLE clause. These will have the number and types
specified in the sampling method’s parameterTypes list, and have been checked to not be null.
seed contains a seed to use for any random numbers generated within the sampling method; it is ei-
ther a hash derived from the REPEATABLE value if one was given, or the result of random() if not.

This function may adjust the fields node->use_bulkread and node->use_pagemode. If
node->use_bulkread is true, which it is by default, the scan will use a buffer access strategy
that encourages recycling buffers after use. It might be reasonable to set this to false if the scan will
visit only a small fraction of the table's pages. If node->use_pagemode is true, which it is by
default, the scan will perform visibility checking in a single pass for all tuples on each visited page. It might be reasonable to set this to false if the scan will select only a small fraction of the tuples on each visited page. That will result in fewer tuple visibility checks being performed, though each one will be more expensive because it will require more locking.

If the sampling method is marked **repeatable_across_scans**, it must be able to select the same set of tuples during a rescan as it did originally, that is a fresh call of BeginSampleScan must lead to selecting the same tuples as before (if the TABLESAMPLE parameters and seed don’t change).

```c
Blockly
NextSampleBlock (SampleScanState *node);
```

Returns the block number of the next page to be scanned, or **InvalidBlockNumber** if no pages remain to be scanned.

This function can be omitted (set the pointer to NULL), in which case the core code will perform a sequential scan of the entire relation. Such a scan can use synchronized scanning, so that the sampling method cannot assume that the relation pages are visited in the same order on each scan.

```c
OffsetNumber
NextSampleTuple (SampleScanState *node,  
    Blockly blockno,  
    OffsetNumber maxoffset);
```

Returns the offset number of the next tuple to be sampled on the specified page, or **InvalidOffsetNumber** if no tuples remain to be sampled. `maxoffset` is the largest offset number in use on the page.

```
Note
NextSampleTuple is not explicitly told which of the offset numbers in the range 1 .. maxoffset actually contain valid tuples. This is not normally a problem since the core code ignores requests to sample missing or invisible tuples; that should not result in any bias in the sample. However, if necessary, the function can examine `node->ss.ss_currentScanDesc->rs_vistuples[]` to identify which tuples are valid and visible. (This requires `node->use_pagemode` to be true.)
```

```
Note
NextSampleTuple must not assume that blockno is the same page number returned by the most recent NextSampleBlock call. It was returned by some previous NextSampleBlock call, but the core code is allowed to call NextSampleBlock in advance of actually scanning pages, so as to support prefetching. It is OK to assume that once sampling of a given page begins, successive NextSampleTuple calls all refer to the same page until InvalidOffsetNumber is returned.
```

```c
void
EndSampleScan (SampleScanState *node);
```

End the scan and release resources. It is normally not important to release palloc’d memory, but any externally-visible resources should be cleaned up. This function can be omitted (set the pointer to NULL) in the common case where no such resources exist.
Chapter 59. Writing A Custom Scan Provider

PostgreSQL supports a set of experimental facilities which are intended to allow extension modules to add new scan types to the system. Unlike a foreign data wrapper, which is only responsible for knowing how to scan its own foreign tables, a custom scan provider can provide an alternative method of scanning any relation in the system. Typically, the motivation for writing a custom scan provider will be to allow the use of some optimization not supported by the core system, such as caching or some form of hardware acceleration. This chapter outlines how to write a new custom scan provider.

Implementing a new type of custom scan is a three-step process. First, during planning, it is necessary to generate access paths representing a scan using the proposed strategy. Second, if one of those access paths is selected by the planner as the optimal strategy for scanning a particular relation, the access path must be converted to a plan. Finally, it must be possible to execute the plan and generate the same results that would have been generated for any other access path targeting the same relation.

59.1. Creating Custom Scan Paths

A custom scan provider will typically add paths for a base relation by setting the following hook, which is called after the core code has generated what it believes to be the complete and correct set of access paths for the relation.

typedef void (*set_rel_pathlist_hook_type) (PlannerInfo *root, RelOptInfo *rel, Index rti, RangeTblEntry *rte);

extern PGDLLIMPORT set_rel_pathlist_hook_type set_rel_pathlist_hook;

Although this hook function can be used to examine, modify, or remove paths generated by the core system, a custom scan provider will typically confine itself to generating CustomPath objects and adding them to rel using add_path. The custom scan provider is responsible for initializing the CustomPath object, which is declared like this:

```c
typedef struct CustomPath
{
    Path      path;
    uint32    flags;
    List     *custom_paths;
    List     *custom_private;
    const CustomPathMethods *methods;
} CustomPath;
```

path must be initialized as for any other path, including the row-count estimate, start and total cost, and sort ordering provided by this path. flags is a bit mask, which should include CUSTOM_PATH_SUPPORT_BACKWARD_SCAN if the custom path can support a backward scan and CUSTOM_PATH_SUPPORT_MARK_RESTORE if it can support mark and restore. Both capabilities are optional. An optional custom_paths is a list of Path nodes used by this custom-path node; these will be transformed into Plan nodes by planner. custom_private can be used to store the custom path's private data. Private data should be stored in a form that can be handled by nodeToString, so that debugging routines that attempt to print the custom path will work as designed. methods must point to a (usually statically allocated) object implementing the required custom path methods, of which there is currently only one. The LibraryName and SymbolName fields must also be initialized so that the dynamic loader can resolve them to locate the method table.
A custom scan provider can also provide join paths. Just as for base relations, such a path must produce the same output as would normally be produced by the join it replaces. To do this, the join provider should set the following hook, and then within the hook function, create `CustomPath` path(s) for the join relation.

```c
typedef void (*set_join_pathlist_hook_type) (PlannerInfo *root,
                                          RelOptInfo *joinrel,
                                          RelOptInfo *outerrel,
                                          RelOptInfo *innerrel,
                                          JoinType jointype,
                                          JoinPathExtraData *extra);
extern PGDLLIMPORT set_join_pathlist_hook_type
    set_join_pathlist_hook;
```

This hook will be invoked repeatedly for the same join relation, with different combinations of inner and outer relations; it is the responsibility of the hook to minimize duplicated work.

### 59.1.1. Custom Scan Path Callbacks

```c
Plan *(*PlanCustomPath) (PlannerInfo *root,
                        RelOptInfo *rel,
                        CustomPath *best_path,
                        List *tlist,
                        List *clauses,
                        List *custom_plans);
```

Convert a custom path to a finished plan. The return value will generally be a `CustomScan` object, which the callback must allocate and initialize. See Section 59.2 for more details.

### 59.2. Creating Custom Scan Plans

A custom scan is represented in a finished plan tree using the following structure:

```c
typedef struct CustomScan
{
    Scan      scan;
    uint32    flags;
    List     *custom_plans;
    List     *custom_exprs;
    List     *custom_private;
    List     *custom_scan_tlist;
    Bitmapset *custom_relids;
    const CustomScanMethods *methods;
} CustomScan;
```

`scan` must be initialized as for any other scan, including estimated costs, target lists, qualifications, and so on. `flags` is a bit mask with the same meaning as in `CustomPath`. `custom_plans` can be used to store child `Plan` nodes. `custom_exprs` should be used to store expression trees that will need to be fixed up by `setrefs.c` and `subselect.c`, while `custom_private` should be used to store other private data that is only used by the custom scan provider itself. `custom_scan_tlist` can be `NIL` when scanning a base relation, indicating that the custom scan returns scan tuples that match the base relation’s row type. Otherwise it is a target list describing the actual scan tuples. `custom_scan_tlist` must be provided for joins, and could be provided for scans if the...
custom scan provider can compute some non-Var expressions. `custom_relids` is set by the core code to the set of relations (range table indexes) that this scan node handles; except when this scan is replacing a join, it will have only one member. `methods` must point to a (usually statically allocated) object implementing the required custom scan methods, which are further detailed below.

When a `CustomScan` scans a single relation, `scan.scanrelid` must be the range table index of the table to be scanned. When it replaces a join, `scan.scanrelid` should be zero.

Plan trees must be able to be duplicated using `copyObject`, so all the data stored within the “custom” fields must consist of nodes that that function can handle. Furthermore, custom scan providers cannot substitute a larger structure that embeds a `CustomScan` for the structure itself, as would be possible for a `CustomPath` or `CustomScanState`.

### 59.2.1. Custom Scan Plan Callbacks

```c
Node *(*CreateCustomScanState) (CustomScan *cscan);
```

Allocate a `CustomScanState` for this `CustomScan`. The actual allocation will often be larger than required for an ordinary `CustomScanState`, because many providers will wish to embed that as the first field of a larger structure. The value returned must have the node tag and `methods` set appropriately, but other fields should be left as zeroes at this stage; after `ExecInitCustomScan` performs basic initialization, the `BeginCustomScan` callback will be invoked to give the custom scan provider a chance to do whatever else is needed.

### 59.3. Executing Custom Scans

When a `CustomScan` is executed, its execution state is represented by a `CustomScanState`, which is declared as follows:

```c
typedef struct CustomScanState
{
    ScanState ss;
    uint32 flags;
    const CustomExecMethods *methods;
} CustomScanState;
```

`ss` is initialized as for any other scan state, except that if the scan is for a join rather than a base relation, `ss.ss_currentRelation` is left `NULL`. `flags` is a bit mask with the same meaning as in `CustomPath` and `CustomScan`. `methods` must point to a (usually statically allocated) object implementing the required custom scan state methods, which are further detailed below. Typically, a `CustomScanState`, which need not support `copyObject`, will actually be a larger structure embedding the above as its first member.

### 59.3.1. Custom Scan Execution Callbacks

```c
void (*BeginCustomScan) (CustomScanState *node,
                         EState *estate,
                         int eflags);
```

Complete initialization of the supplied `CustomScanState`. Standard fields have been initialized by `ExecInitCustomScan`, but any private fields should be initialized here.

```c
TupleTableSlot *(*ExecCustomScan) (CustomScanState *node);
```
Fetch the next scan tuple. If any tuples remain, it should fill ps_ResultTupleSlot with the next tuple in the current scan direction, and then return the tuple slot. If not, NULL or an empty slot should be returned.

```c
void (*EndCustomScan) (CustomScanState *node);
```

Clean up any private data associated with the CustomScanState. This method is required, but it does not need to do anything if there is no associated data or it will be cleaned up automatically.

```c
void (*ReScanCustomScan) (CustomScanState *node);
```

Rewind the current scan to the beginning and prepare to rescan the relation.

```c
void (*MarkPosCustomScan) (CustomScanState *node);
```

Save the current scan position so that it can subsequently be restored by the RestrPosCustomScan callback. This callback is optional, and need only be supplied if the CUSTOMPATH_SUP-PORT_MARK_RESTORE flag is set.

```c
void (*RestrPosCustomScan) (CustomScanState *node);
```

Restore the previous scan position as saved by the MarkPosCustomScan callback. This callback is optional, and need only be supplied if the CUSTOMPATH_SUPPORT_MARK_RESTORE flag is set.

```c
Size (*EstimateDSMCustomScan) (CustomScanState *node, ParallelContext *pcxt);
```

Estimate the amount of dynamic shared memory that will be required for parallel operation. This may be higher than the amount that will actually be used, but it must not be lower. The return value is in bytes. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution.

```c
void (*InitializeDSMCustomScan) (CustomScanState *node, ParallelContext *pcxt, void *coordinate);
```

Initialize the dynamic shared memory that will be required for parallel operation. coordinate points to a shared memory area of size equal to the return value of EstimateDSMCustomScan. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution.

```c
void (*ReInitializeDSMCustomScan) (CustomScanState *node, ParallelContext *pcxt, void *coordinate);
```

Re-initialize the dynamic shared memory required for parallel operation when the custom-scan plan node is about to be re-scanned. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution. Recommended practice is that this callback reset only shared state, while the ReScanCustomScan callback resets only local state. Currently, this callback will be called before ReScanCustomScan, but it’s best not to rely on that ordering.

```c
void (*InitializeWorkerCustomScan) (CustomScanState *node, shm_toc *toc, void *coordinate);
```
Initialize a parallel worker's local state based on the shared state set up by the leader during `InitializeDSMCustomScan`. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution.

```c
void (*ShutdownCustomScan) (CustomScanState *node);
```

Release resources when it is anticipated the node will not be executed to completion. This is not called in all cases; sometimes, `EndCustomScan` may be called without this function having been called first. Since the DSM segment used by parallel query is destroyed just after this callback is invoked, custom scan providers that wish to take some action before the DSM segment goes away should implement this method.

```c
void (*ExplainCustomScan) (CustomScanState *node,
                           List *ancestors,
                           ExplainState *es);
```

Output additional information for `EXPLAIN` of a custom-scan plan node. This callback is optional. Common data stored in the `ScanState`, such as the target list and scan relation, will be shown even without this callback, but the callback allows the display of additional, private state.
Chapter 60. Genetic Query Optimizer

60.1. Query Handling as a Complex Optimization Problem

Among all relational operators the most difficult one to process and optimize is the *join*. The number of possible query plans grows exponentially with the number of joins in the query. Further optimization effort is caused by the support of a variety of *join methods* (e.g., nested loop, hash join, merge join in PostgreSQL) to process individual joins and a diversity of *indexes* (e.g., B-tree, hash, GiST and GIN in PostgreSQL) as access paths for relations.

The normal PostgreSQL query optimizer performs a *near-exhaustive search* over the space of alternative strategies. This algorithm, first introduced in IBM's System R database, produces a near-optimal join order, but can take an enormous amount of time and memory space when the number of joins in the query grows large. This makes the ordinary PostgreSQL query optimizer inappropriate for queries that join a large number of tables.

The Institute of Automatic Control at the University of Mining and Technology, in Freiberg, Germany, encountered some problems when it wanted to use PostgreSQL as the backend for a decision support knowledge based system for the maintenance of an electrical power grid. The DBMS needed to handle large join queries for the inference machine of the knowledge based system. The number of joins in these queries made using the normal query optimizer infeasible.

In the following we describe the implementation of a *genetic algorithm* to solve the join ordering problem in a manner that is efficient for queries involving large numbers of joins.

60.2. Genetic Algorithms

The genetic algorithm (GA) is a heuristic optimization method which operates through randomized search. The set of possible solutions for the optimization problem is considered as a *population* of *individuals*. The degree of adaptation of an individual to its environment is specified by its *fitness*.

The coordinates of an individual in the search space are represented by *chromosomes*, in essence a set of character strings. A *gene* is a subsection of a chromosome which encodes the value of a single parameter being optimized. Typical encodings for a gene could be *binary* or *integer*.

Through simulation of the evolutionary operations *recombination*, *mutation*, and *selection* new generations of search points are found that show a higher average fitness than their ancestors.

According to the comp.ai.genetic FAQ it cannot be stressed too strongly that a GA is not a pure random search for a solution to a problem. A GA uses stochastic processes, but the result is distinctly non-random (better than random).
Genetic Query Optimizer

Figure 60.1. Structured Diagram of a Genetic Algorithm

\[ P(t) \] \text{generation of ancestors at a time } t \\
\[ P''(t) \] \text{generation of descendants at a time } t 

\begin{verbatim}
+=========================================+
| >>>>>>>>>> Algorithm GA <<<<<<<<<<<+
+=========================================+
| INITIALIZE t := 0                      |
+=========================================+
| INITIALIZE P(t)                        |
+=========================================+
| evaluate FITNESS of P(t)              |
+=========================================+
| while not STOPPING CRITERION do       |
+=========================================+
|   | P'(t) := RECOMBINATION{P(t)}        |
+=========================================+
|   | P''(t) := MUTATION{P'(t)}           |
+=========================================+
|   | P(t+1) := SELECTION{P''(t) + P(t)}  |
+=========================================+
|   | evaluate FITNESS of P''(t)         |
+=========================================+
|   | t := t + 1                         |
+=========================================+
\end{verbatim}

60.3. Genetic Query Optimization (GEQO) in PostgreSQL

The GEQO module approaches the query optimization problem as though it were the well-known traveling salesman problem (TSP). Possible query plans are encoded as integer strings. Each string represents the join order from one relation of the query to the next. For example, the join tree

\[
\begin{array}{c}
/ \\
/ 2 \\
/ \ 3 \\
4 1
\end{array}
\]

is encoded by the integer string ‘4-1-3-2’, which means, first join relation ‘4’ and ‘1’, then ‘3’, and then ‘2’, where 1, 2, 3, 4 are relation IDs within the PostgreSQL optimizer.

Specific characteristics of the GEQO implementation in PostgreSQL are:

- Usage of a \textit{steady state} GA (replacement of the least fit individuals in a population, not whole-generational replacement) allows fast convergence towards improved query plans. This is essential for query handling with reasonable time;
- Usage of \textit{edge recombination crossover} which is especially suited to keep edge losses low for the solution of the TSP by means of a GA;
- Mutation as genetic operator is deprecated so that no repair mechanisms are needed to generate legal TSP tours.

Parts of the GEQO module are adapted from D. Whitley’s Genitor algorithm.
The GEQO module allows the PostgreSQL query optimizer to support large join queries effectively through non-exhaustive search.

60.3.1. Generating Possible Plans with GEQO

The GEQO planning process uses the standard planner code to generate plans for scans of individual relations. Then join plans are developed using the genetic approach. As shown above, each candidate join plan is represented by a sequence in which to join the base relations. In the initial stage, the GEQO code simply generates some possible join sequences at random. For each join sequence considered, the standard planner code is invoked to estimate the cost of performing the query using that join sequence. (For each step of the join sequence, all three possible join strategies are considered; and all the initially-determined relation scan plans are available. The estimated cost is the cheapest of these possibilities.) Join sequences with lower estimated cost are considered “more fit” than those with higher cost. The genetic algorithm discards the least fit candidates. Then new candidates are generated by combining genes of more-fit candidates — that is, by using randomly-chosen portions of known low-cost join sequences to create new sequences for consideration. This process is repeated until a preset number of join sequences have been considered; then the best one found at any time during the search is used to generate the finished plan.

This process is inherently nondeterministic, because of the randomized choices made during both the initial population selection and subsequent “mutation” of the best candidates. To avoid surprising changes of the selected plan, each run of the GEQO algorithm restarts its random number generator with the current `geqo_seed` parameter setting. As long as `geqo_seed` and the other GEQO parameters are kept fixed, the same plan will be generated for a given query (and other planner inputs such as statistics). To experiment with different search paths, try changing `geqo_seed`.

60.3.2. Future Implementation Tasks for PostgreSQL GEQO

Work is still needed to improve the genetic algorithm parameter settings. In file `src/backend/optimizer/geqo/geqo_main.c`, routines `gimme_pool_size` and `gimme_number_generations`, we have to find a compromise for the parameter settings to satisfy two competing demands:

- Optimality of the query plan
- Computing time

In the current implementation, the fitness of each candidate join sequence is estimated by running the standard planner’s join selection and cost estimation code from scratch. To the extent that different candidates use similar sub-sequences of joins, a great deal of work will be repeated. This could be made significantly faster by retaining cost estimates for sub-joins. The problem is to avoid expending unreasonable amounts of memory on retaining that state.

At a more basic level, it is not clear that solving query optimization with a GA algorithm designed for TSP is appropriate. In the TSP case, the cost associated with any substring (partial tour) is independent of the rest of the tour, but this is certainly not true for query optimization. Thus it is questionable whether edge recombination crossover is the most effective mutation procedure.

60.4. Further Reading

The following resources contain additional information about genetic algorithms:

- Evolutionary Computation and its application to art and design², by Craig Reynolds

¹ http://www.aip.de/~ast/EvolCompFAQ/
² http://www.red3d.com/cwr/evolve.html
Genetic Query Optimizer

- [elma04]
- [fong]
Chapter 61. Index Access Method Interface Definition

This chapter defines the interface between the core PostgreSQL system and index access methods, which manage individual index types. The core system knows nothing about indexes beyond what is specified here, so it is possible to develop entirely new index types by writing add-on code.

All indexes in PostgreSQL are what are known technically as secondary indexes; that is, the index is physically separate from the table file that it describes. Each index is stored as its own physical relation and so is described by an entry in the pg_class catalog. The contents of an index are entirely under the control of its index access method. In practice, all index access methods divide indexes into standard-size pages so that they can use the regular storage manager and buffer manager to access the index contents. (All the existing index access methods furthermore use the standard page layout described in Section 68.6, and most use the same format for index tuple headers; but these decisions are not forced on an access method.)

An index is effectively a mapping from some data key values to tuple identifiers, or TIDs, of row versions (tuples) in the index's parent table. A TID consists of a block number and an item number within that block (see Section 68.6). This is sufficient information to fetch a particular row version from the table. Indexes are not directly aware that under MVCC, there might be multiple extant versions of the same logical row; to an index, each tuple is an independent object that needs its own index entry. Thus, an update of a row always creates all-new index entries for the row, even if the key values did not change. (HOT tuples are an exception to this statement; but indexes do not deal with those, either.) Index entries for dead tuples are reclaimed (by vacuuming) when the dead tuples themselves are reclaimed.

61.1. Basic API Structure for Indexes

Each index access method is described by a row in the pg_am system catalog. The pg_am entry specifies a name and a handler function for the access method. These entries can be created and deleted using the CREATE ACCESS METHOD and DROP ACCESS METHOD SQL commands.

An index access method handler function must be declared to accept a single argument of type internal and to return the pseudo-type index_am_handler. The argument is a dummy value that simply serves to prevent handler functions from being called directly from SQL commands. The result of the function must be a palloc'd struct of type IndexAmRoutine, which contains everything that the core code needs to know to make use of the index access method. The IndexAmRoutine struct, also called the access method's API struct, includes fields specifying assorted fixed properties of the access method, such as whether it can support multicolumn indexes. More importantly, it contains pointers to support functions for the access method, which do all of the real work to access indexes. These support functions are plain C functions and are not visible or callable at the SQL level. The support functions are described in Section 61.2.

The structure IndexAmRoutine is defined thus:

typedef struct IndexAmRoutine
{
    NodeTag     type;

    /*
     * Total number of strategies (operators) by which we can traverse/search
     * this AM.  Zero if AM does not have a fixed set of strategy assignments.
     */
}
uint16  amstrategies;
/* total number of support functions that this AM uses */
uint16  amsupport;
/* does AM support ORDER BY indexed column's value? */
bool    amcanorder;
/* does AM support ORDER BY result of an operator on indexed column? */
bool    amcanorderbyop;
/* does AM support backward scanning? */
bool    amcanbackward;
/* does AM support UNIQUE indexes? */
bool    amcanunique;
/* does AM support multi-column indexes? */
bool    amcanmulticol;
/* does AM require scans to have a constraint on the first index column? */
bool    amoptionalkey;
/* does AM handle ScalarArrayOpExpr quals? */
bool    amsearcharray;
/* does AM handle IS NULL/IS NOT NULL quals? */
bool    amsearchnulls;
/* can index storage data type differ from column data type? */
bool    amstorage;
/* can an index of this type be clustered on? */
bool    amclusterable;
/* does AM handle predicate locks? */
bool    ampredlocks;
/* does AM support parallel scan? */
bool    amcanparallel;
/* does AM support columns included with clause INCLUDE? */
bool    amcaninclude;
/* type of data stored in index, or InvalidOid if variable */
Oid     amkeytype;

/* interface functions */
ambuild_function ambuild;
ambuildempty_function ambuildempty;
aminsert_function aminsert;
ambulkdelete_function ambulkdelete;
amvacuumcleanup_function amvacuumcleanup;
amcanreturn_function amcanreturn; /* can be NULL */
amcostestimate_function amcostestimate;
amoptions_function amoptions;
amproperty_function amproperty; /* can be NULL */
amvalidate_function amvalidate;
ambeginscan_function ambeginscan;
amrescan_function amrescan;
amgettuple_function amgettuples; /* can be NULL */
amgetbitmap_function amgetbitmap; /* can be NULL */
amendscan_function amendscan;
ammarkpos_function ammarkpos; /* can be NULL */
amrestrpos_function amrestrpos; /* can be NULL */

/* interface functions to support parallel index scans */
amestimateparallelscan_function amestimateparallelscan; /* can be NULL */
aminitparallelscan_function aminitparallelscan; /* can be NULL */
To be useful, an index access method must also have one or more operator families and operator
classes defined in `pg_opfamily`, `pg_opclass`, `pg_amop`, and `pg_amproc`. These entries allow
the planner to determine what kinds of query qualifications can be used with indexes of this access
method. Operator families and classes are described in Section 38.15, which is prerequisite material
for reading this chapter.

An individual index is defined by a `pg_class` entry that describes it as a physical relation, plus a
`pg_index` entry that shows the logical content of the index — that is, the set of index columns it
has and the semantics of those columns, as captured by the associated operator classes. The index
columns (key values) can be either simple columns of the underlying table or expressions over the
table rows. The index access method normally has no interest in where the index key values come
from (it is always handed precomputed key values) but it will be very interested in the operator class
information in `pg_index`. Both of these catalog entries can be accessed as part of the Relation
data structure that is passed to all operations on the index.

Some of the flag fields of `IndexAmRoutine` have nonobvious implications. The requirements of
`amcanunique` are discussed in Section 61.5. The `amcanmulticol` flag asserts that the access
method supports multicolumn indexes, while `amoptionalkey` asserts that it allows scans where
no indexable restriction clause is given for the first index column. When `amcanmulticol` is false,
`amoptionalkey` essentially says whether the access method supports full-index scans without any
restriction clause. Access methods that support multiple index columns must support scans that omit
restrictions on any or all of the columns after the first; however they are permitted to require some
restriction to appear for the first index column, and this is signaled by setting `amoptionalkey` false.
One reason that an index AM might set `amoptionalkey` false is if it doesn't index null values.
Since most indexable operators are strict and hence cannot return true for null inputs, it is at first sight
attractive to not store index entries for null values: they could never be returned by an index scan
anyway. However, this argument fails when an index scan has no restriction clause for a given index
column. In practice this means that indexes that have `amoptionalkey` true must index nulls, since
the planner might decide to use such an index with no scan keys at all. A related restriction is that
an index access method that supports multiple index columns must support indexing null values in
columns after the first, because the planner will assume the index can be used for queries that do not
restrict these columns. For example, consider an index on (a,b) and a query with `WHERE a = 4`.
The system will assume the index can be used to scan for rows with `a = 4`, which is wrong if the
index omits rows where `b` is null. It is, however, OK to omit rows where the first indexed column is
null. An index access method that does index nulls may also set `amsearchnulls`, indicating that it
supports `IS NULL` and `IS NOT NULL` clauses as search conditions.

## 61.2. Index Access Method Functions

The index construction and maintenance functions that an index access method must provide in `IndexAmRoutine` are:

```c
IndexBuildResult *
ambuild (Relation heapRelation,
      Relation indexRelation,
      IndexInfo *indexInfo);
```

Build a new index. The index relation has been physically created, but is empty. It must be filled in
with whatever fixed data the access method requires, plus entries for all tuples already existing in the
table. Ordinarily the `ambuild` function will call `IndexBuildHeapScan()` to scan the table for
existing tuples and compute the keys that need to be inserted into the index. The function must return
a `palloc`d struct containing statistics about the new index.
void
ambuildempty (Relation indexRelation);

Build an empty index, and write it to the initialization fork (INIT_FORKNUM) of the given relation. This method is called only for unlogged indexes; the empty index written to the initialization fork will be copied over the main relation fork on each server restart.

bool
aminsert (Relation indexRelation,
     Datum *values,
     bool *isnull,
     ItemPointer heap_tid,
     Relation heapRelation,
     IndexUniqueCheck checkUnique,
     IndexInfo *indexInfo);

Insert a new tuple into an existing index. The values and isnull arrays give the key values to be indexed, and heap_tid is the TID to be indexed. If the access method supports unique indexes (its amcanunique flag is true) then checkUnique indicates the type of uniqueness check to perform. This varies depending on whether the unique constraint is deferrable; see Section 61.5 for details. Normally the access method only needs the heapRelation parameter when performing uniqueness checking (since then it will have to look into the heap to verify tuple liveness).

The function's Boolean result value is significant only when checkUnique is UNIQUE_CHECK_PARTIAL. In this case a true result means the new entry is known unique, whereas false means it might be non-unique (and a deferred uniqueness check must be scheduled). For other cases a constant false result is recommended.

Some indexes might not index all tuples. If the tuple is not to be indexed, aminsert should just return without doing anything.

If the index AM wishes to cache data across successive index insertions within a SQL statement, it can allocate space in indexInfo->ii_Context and store a pointer to the data in indexInfo->ii_AmCache (which will be NULL initially).

IndexBulkDeleteResult *
ambulkdelete (IndexVacuumInfo *info,
        IndexBulkDeleteResult *stats,
        IndexBulkDeleteCallback callback,
        void *callback_state);

Delete tuple(s) from the index. This is a “bulk delete” operation that is intended to be implemented by scanning the whole index and checking each entry to see if it should be deleted. The passed-in callback function must be called, in the style callback(TID, callback_state) returns bool, to determine whether any particular index entry, as identified by its referenced TID, is to be deleted. Must return either NULL or a palloc’d struct containing statistics about the effects of the deletion operation. It is OK to return NULL if no information needs to be passed on to amvacuumcleanup.

Because of limited maintenance_work_mem, ambulkdelete might need to be called more than once when many tuples are to be deleted. The stats argument is the result of the previous call for this index (it is NULL for the first call within a VACUUM operation). This allows the AM to accumulate statistics across the whole operation. Typically, ambulkdelete will modify and return the same struct if the passed stats is not null.

IndexBulkDeleteResult *
amvacuumcleanup (IndexVacuumInfo *info,  
    IndexBulkDeleteResult *stats);

Clean up after a VACUUM operation (zero or more ambulkdelete calls). This does not have to do anything beyond returning index statistics, but it might perform bulk cleanup such as reclaiming empty index pages. stats is whatever the last ambulkdelete call returned, or NULL if ambulkdelete was not called because no tuples needed to be deleted. If the result is not NULL it must be a palloc'd struct. The statistics it contains will be used to update pg_class, and will be reported by VACUUM if VERBOSE is given. It is OK to return NULL if the index was not changed at all during the VACUUM operation, but otherwise correct stats should be returned.

As of PostgreSQL 8.4, amvacuumcleanup will also be called at completion of an ANALYZE operation. In this case stats is always NULL and any return value will be ignored. This case can be distinguished by checking info->analyze_only. It is recommended that the access method do nothing except post-insert cleanup in such a call, and that only in an autovacuum worker process.

bool
amcanreturn (Relation indexRelation, int attno);

Check whether the index can support index-only scans on the given column, by returning the indexed column values for an index entry in the form of an IndexTuple. The attribute number is 1-based, i.e. the first column's attno is 1. Returns true if supported, else false. If the access method does not support index-only scans at all, the amcanreturn field in its IndexAmRoutine struct can be set to NULL.

void
amcostestimate (PlannerInfo *root,  
    IndexPath *path,  
    double loop_count,  
    Cost *indexStartupCost,  
    Cost *indexTotalCost,  
    Selectivity *indexSelectivity,  
    double *indexCorrelation,  
    double *indexPages);

Estimate the costs of an index scan. This function is described fully in Section 61.6, below.

bytea *
amoptions (ArrayType *reloptions,  
             bool validate);

Parse and validate the reoptions array for an index. This is called only when a non-null reoptions array exists for the index. reoptions is a text array containing entries of the form name=value. The function should construct a bytea value, which will be copied into the rd_options field of the index's relcache entry. The data contents of the bytea value are open for the access method to define; most of the standard access methods use struct StdRdOptions. When validate is true, the function should report a suitable error message if any of the options are unrecognized or have invalid values; when validate is false, invalid entries should be silently ignored. (validate is false when loading options already stored in pg_catalog; an invalid entry could only be found if the access method has changed its rules for options, and in that case ignoring obsolete entries is appropriate.) It is OK to return NULL if default behavior is wanted.

bool
amproperty (Oid index_oid, int attno,  
             IndexAMProperty prop, const char *propname,  
             bool *res, bool *isnull);
The amproperty method allows index access methods to override the default behavior of pg_index_column_has_property and related functions. If the access method does not have any special behavior for index property inquiries, the amproperty field in its IndexAmRoutine struct can be set to NULL. Otherwise, the amproperty method will be called with index_oid and attno both zero for pg_indexam_has_property calls, or with index_oid valid and attno zero for pg_index_has_property calls, or with index_oid valid and attno greater than zero for pg_index_column_has_property calls. prop is an enum value identifying the property being tested, while propname is the original property name string. If the core code does not recognize the property name then prop is AMPROP_UNKNOWN. Access methods can define custom property names by checking propname for a match (use pg_strcasecmp to match, for consistency with the core code); for names known to the core code, it’s better to inspect prop. If the amproperty method returns true then it has determined the property test result: it must set *res to the boolean value to return, or set *isnull to true to return a NULL. (Both of the referenced variables are initialized to false before the call.) If the amproperty method returns false then the core code will proceed with its normal logic for determining the property test result.

Access methods that support ordering operators should implement AMPROP_DISTANCE_ORDERABLE property testing, as the core code does not know how to do that and will return NULL. It may also be advantageous to implement AMPROP_RETURNABLE testing, if that can be done more cheaply than by opening the index and calling amcanreturn, which is the core code’s default behavior. The default behavior should be satisfactory for all other standard properties.

```c
bool amvalidate (Oid opclassoid);
```

Validate the catalog entries for the specified operator class, so far as the access method can reasonably do that. For example, this might include testing that all required support functions are provided. The amvalidate function must return false if the opclass is invalid. Problems should be reported with ereport messages.

The purpose of an index, of course, is to support scans for tuples matching an indexable WHERE condition, often called a qualifier or scan key. The semantics of index scanning are described more fully in Section 61.3, below. An index access method can support “plain” index scans, “bitmap” index scans, or both. The scan-related functions that an index access method must or may provide are:

```c
IndexScanDesc
ambeginscan (Relation indexRelation,
            int nkeys,
            int norderbys);
```

Prepare for an index scan. The nkeys and norderbys parameters indicate the number of quals and ordering operators that will be used in the scan; these may be useful for space allocation purposes. Note that the actual values of the scan keys aren’t provided yet. The result must be a palloc’d struct. For implementation reasons the index access method must create this struct by calling RelationGetIndexScan(). In most cases ambeginscan does little beyond making that call and perhaps acquiring locks; the interesting parts of index-scan startup are in amrescan.

```c
void
amrescan (IndexScanDesc scan,
            ScanKey keys,
            int nkeys,
            ScanKey orderbys,
            int norderbys);
```

Start or restart an index scan, possibly with new scan keys. (To restart using previously-passed keys, NULL is passed for keys and/or orderbys.) Note that it is not allowed for the number of keys or order-by operators to be larger than what was passed to ambeginscan. In practice the restart feature
Index Access Method
Interface Definition

is used when a new outer tuple is selected by a nested-loop join and so a new key comparison value is needed, but the scan key structure remains the same.

```c
boolean
amgettuple (IndexScanDesc scan,
            ScanDirection direction);
```

Fetch the next tuple in the given scan, moving in the given direction (forward or backward in the index). Returns true if a tuple was obtained, false if no matching tuples remain. In the true case the tuple TID is stored into the scan structure. Note that “success” means only that the index contains an entry that matches the scan keys, not that the tuple necessarily still exists in the heap or will pass the caller's snapshot test. On success, amgettuple must also set scan->xs_recheck to true or false. False means it is certain that the index entry matches the scan keys. true means this is not certain, and the conditions represented by the scan keys must be rechecked against the heap tuple after fetching it. This provision supports “lossy” index operators. Note that rechecking will extend only to the scan conditions; a partial index predicate (if any) is never rechecked by amgettuple callers.

If the index supports index-only scans (i.e., amcanreturn returns true for it), then on success the AM must also check scan->xs_want_itup, and if that is true it must return the originally indexed data for the index entry. The data can be returned in the form of an IndexTuple pointer stored at scan->xs_itup, with tuple descriptor scan->xs_itupdesc; or in the form of a HeapTuple pointer stored at scan->xs_hitup, with tuple descriptor scan->xs_hitupdesc. (The latter format should be used when reconstructing data that might possibly not fit into an IndexTuple.) In either case, management of the data referenced by the pointer is the access method's responsibility. The data must remain good at least until the next amgettuple, amrescan, or amendscan call for the scan.

The amgettuple function need only be provided if the access method supports “plain” index scans. If it doesn't, the amgettuple field in its IndexAmRoutine struct must be set to NULL.

```c
int64
amgetbitmap (IndexScanDesc scan,
              TIDBitmap *tbm);
```

Fetch all tuples in the given scan and add them to the caller-supplied TIDBitmap (that is, OR the set of tuple IDs into whatever set is already in the bitmap). The number of tuples fetched is returned (this might be just an approximate count, for instance some AMs do not detect duplicates). While inserting tuple IDs into the bitmap, amgetbitmap can indicate that rechecking of the scan conditions is required for specific tuple IDs. This is analogous to the xs_recheck output parameter of amgettuple. Note: in the current implementation, support for this feature is conflated with support for lossy storage of the bitmap itself, and therefore callers recheck both the scan conditions and the partial index predicate (if any) for recheckable tuples. That might not always be true, however. amgetbitmap and amgettuple cannot be used in the same index scan; there are other restrictions too when using amgetbitmap, as explained in Section 61.3.

The amgetbitmap function need only be provided if the access method supports “bitmap” index scans. If it doesn't, the amgetbitmap field in its IndexAmRoutine struct must be set to NULL.

```c
void
amendscan (IndexScanDesc scan);
```

End a scan and release resources. The scan struct itself should not be freed, but any locks or pins taken internally by the access method must be released, as well as any other memory allocated by ambeginscan and other scan-related functions.
**Index Scanning**

In an index scan, the index access method is responsible for regurgitating the TIDs of all the tuples it has been told about that match the *scan keys*. The access method is *not* involved in actually fetch-
ing those tuples from the index's parent table, nor in determining whether they pass the scan's time qualification test or other conditions.

A scan key is the internal representation of a WHERE clause of the form \textit{index\_key \text{operator} constant}, where the index key is one of the columns of the index and the operator is one of the members of the operator family associated with that index column. An index scan has zero or more scan keys, which are implicitly ANDed — the returned tuples are expected to satisfy all the indicated conditions.

The access method can report that the index is \textit{lossy}, or requires rechecks, for a particular query. This implies that the index scan will return all the entries that pass the scan key, plus possibly additional entries that do not. The core system's index-scan machinery will then apply the index conditions again to the heap tuple to verify whether or not it really should be selected. If the recheck option is not specified, the index scan must return exactly the set of matching entries.

Note that it is entirely up to the access method to ensure that it correctly finds all and only the entries passing all the given scan keys. Also, the core system will simply hand off all the WHERE clauses that match the index keys and operator families, without any semantic analysis to determine whether they are redundant or contradictory. As an example, given \textit{WHERE x > 4 AND x > 14} where \textit{x} is a b-tree indexed column, it is left to the b-tree \textit{amrescan} function to realize that the first scan key is redundant and can be discarded. The extent of preprocessing needed during \textit{amrescan} will depend on the extent to which the index access method needs to reduce the scan keys to a "normalized" form.

Some access methods return index entries in a well-defined order, others do not. There are actually two different ways that an access method can support sorted output:

- Access methods that always return entries in the natural ordering of their data (such as btree) should set \textit{amcanorder} to true. Currently, such access methods must use btree-compatible strategy numbers for their equality and ordering operators.

- Access methods that support ordering operators should set \textit{amcanorderbyop} to true. This indicates that the index is capable of returning entries in an order satisfying \textit{ORDER BY index\_key \text{operator} constant}. Scan modifiers of that form can be passed to \textit{amrescan} as described previously.

The \textit{amgettuple} function has a \textit{direction} argument, which can be either \textit{ForwardScanDirection} (the normal case) or \textit{BackwardScanDirection}. If the first call after \textit{amrescan} specifies \textit{BackwardScanDirection}, then the set of matching index entries is to be scanned back-to-front rather than in the normal front-to-back direction, so \textit{amgettuple} must return the last matching tuple in the index, rather than the first one as it normally would. (This will only occur for access methods that set \textit{amcanorder} to true.) After the first call, \textit{amgettuple} must be prepared to advance the scan in either direction from the most recently returned entry. (But if \textit{amcanbackward} is false, all subsequent calls will have the same direction as the first one.)

Access methods that support ordered scans must support "marking" a position in a scan and later returning to the marked position. The same position might be restored multiple times. However, only one position need be remembered per scan; a new \textit{ammarkpos} call overrides the previously marked position. An access method that does not support ordered scans need not provide \textit{ammarkpos} and \textit{amrestrpos} functions in \textit{IndexAmRoutine}; set those pointers to NULL instead.

Both the scan position and the mark position (if any) must be maintained consistently in the face of concurrent insertions or deletions in the index. It is OK if a freshly-inserted entry is not returned by a scan that would have found the entry if it had existed when the scan started, or for the scan to return such an entry upon rescanning or backing up even though it had not been returned the first time through. Similarly, a concurrent delete might or might not be reflected in the results of a scan. What is important is that insertions or deletions not cause the scan to miss or multiply return entries that were not themselves being inserted or deleted.

If the index stores the original indexed data values (and not some lossy representation of them), it is useful to support index-only scans, in which the index returns the actual data not just the TID of
the heap tuple. This will only avoid I/O if the visibility map shows that the TID is on an all-visible page; else the heap tuple must be visited anyway to check MVCC visibility. But that is no concern of the access method’s.

Instead of using amgettuple, an index scan can be done with amgetbitmap to fetch all tuples in one call. This can be noticeably more efficient than amgettuple because it allows avoiding lock/unlock cycles within the access method. In principle amgetbitmap should have the same effects as repeated amgettuple calls, but we impose several restrictions to simplify matters. First of all, amgetbitmap returns all tuples at once and marking or restoring scan positions isn’t supported. Secondly, the tuples are returned in a bitmap which doesn’t have any specific ordering, which is why amgetbitmap doesn’t take a direction argument. (Ordering operators will never be supplied for such a scan, either.) Also, there is no provision for index-only scans with amgetbitmap, since there is no way to return the contents of index tuples. Finally, amgetbitmap does not guarantee any locking of the returned tuples, with implications spelled out in Section 61.4.

Note that it is permitted for an access method to implement only amgetbitmap and not amgettuple, or vice versa, if its internal implementation is unsuited to one API or the other.

61.4. Index Locking Considerations

Index access methods must handle concurrent updates of the index by multiple processes. The core PostgreSQL system obtains AccessShareLock on the index during an index scan, and RowExclusiveLock when updating the index (including plain VACUUM). Since these lock types do not conflict, the access method is responsible for handling any fine-grained locking it might need. An exclusive lock on the index as a whole will be taken only during index creation, destruction, or REINDEX.

Building an index type that supports concurrent updates usually requires extensive and subtle analysis of the required behavior. For the b-tree and hash index types, you can read about the design decisions involved in src/backend/access/nbtree/README and src/backend/access/hash/README.

Aside from the index’s own internal consistency requirements, concurrent updates create issues about consistency between the parent table (the heap) and the index. Because PostgreSQL separates accesses and updates of the heap from those of the index, there are windows in which the index might be inconsistent with the heap. We handle this problem with the following rules:

- A new heap entry is made before making its index entries. (Therefore a concurrent index scan is likely to fail to see the heap entry. This is okay because the index reader would be uninterested in an uncommitted row anyway. But see Section 61.5.)

- When a heap entry is to be deleted (by VACUUM), all its index entries must be removed first.

- An index scan must maintain a pin on the index page holding the item last returned by amgettuple, and ambulkdelete cannot delete entries from pages that are pinned by other backends. The need for this rule is explained below.

Without the third rule, it is possible for an index reader to see an index entry just before it is removed by VACUUM, and then to arrive at the corresponding heap entry after that was removed by VACUUM. This creates no serious problems if that item number is still unused when the reader reaches it, since an empty item slot will be ignored by heap_fetch(). But what if a third backend has already reused the item slot for something else? When using an MVCC-compliant snapshot, there is no problem because the new occupant of the slot is certain to be too new to pass the snapshot test. However, with a non-MVCC-compliant snapshot (such as SnapshotAny), it would be possible to accept and return a row that does not in fact match the scan keys. We could defend against this scenario by requiring the scan keys to be rechecked against the heap row in all cases, but that is too expensive. Instead, we use a pin on an index page as a proxy to indicate that the reader might still be “in flight” from the index entry to the matching heap entry. Making ambulkdelete block on such a pin ensures that VACUUM cannot delete the heap entry before the reader is done with it. This solution costs little in run time, and adds blocking overhead only in the rare cases where there actually is a conflict.
This solution requires that index scans be "synchronous": we have to fetch each heap tuple immediately after scanning the corresponding index entry. This is expensive for a number of reasons. An "asynchronous" scan in which we collect many TIDs from the index, and only visit the heap tuples sometime later, requires much less index locking overhead and can allow a more efficient heap access pattern. Per the above analysis, we must use the synchronous approach for non-MVCC-compliant snapshots, but an asynchronous scan is workable for a query using an MVCC snapshot.

In an amgetbitmap index scan, the access method does not keep an index pin on any of the returned tuples. Therefore it is only safe to use such scans with MVCC-compliant snapshots.

When the ampredlocks flag is not set, any scan using that index access method within a serializable transaction will acquire a nonblocking predicate lock on the full index. This will generate a read-write conflict with the insert of any tuple into that index by a concurrent serializable transaction. If certain patterns of read-write conflicts are detected among a set of concurrent serializable transactions, one of those transactions may be canceled to protect data integrity. When the flag is set, it indicates that the index access method implements finer-grained predicate locking, which will tend to reduce the frequency of such transaction cancellations.

61.5. Index Uniqueness Checks

PostgreSQL enforces SQL uniqueness constraints using unique indexes, which are indexes that disallow multiple entries with identical keys. An access method that supports this feature sets amcanunique true. (At present, only b-tree supports it.) Columns listed in the INCLUDE clause are not considered when enforcing uniqueness.

Because of MVCC, it is always necessary to allow duplicate entries to exist physically in an index: the entries might refer to successive versions of a single logical row. The behavior we actually want to enforce is that no MVCC snapshot could include two rows with equal index keys. This breaks down into the following cases that must be checked when inserting a new row into a unique index:

• If a conflicting valid row has been deleted by the current transaction, it's okay. (In particular, since an UPDATE always deletes the old row version before inserting the new version, this will allow an UPDATE on a row without changing the key.)

• If a conflicting row has been inserted by an as-yet-uncommitted transaction, the would-be inserter must wait to see if that transaction commits. If it rolls back then there is no conflict. If it commits without deleting the conflicting row again, there is a uniqueness violation. (In practice we just wait for the other transaction to end and then redo the visibility check in toto.)

• Similarly, if a conflicting valid row has been deleted by an as-yet-uncommitted transaction, the would-be inserter must wait for that transaction to commit or abort, and then repeat the test.

Furthermore, immediately before reporting a uniqueness violation according to the above rules, the access method must recheck the liveness of the row being inserted. If it is committed dead then no violation should be reported. (This case cannot occur during the ordinary scenario of inserting a row that's just been created by the current transaction. It can happen during CREATE UNIQUE INDEX CONCURRENTLY, however.)

We require the index access method to apply these tests itself, which means that it must reach into the heap to check the commit status of any row that is shown to have a duplicate key according to the index contents. This is without a doubt ugly and non-modular, but it saves redundant work: if we did a separate probe then the index lookup for a conflicting row would be essentially repeated while finding the place to insert the new row's index entry. What's more, there is no obvious way to avoid race conditions unless the conflict check is an integral part of insertion of the new index entry.

If the unique constraint is deferrable, there is additional complexity: we need to be able to insert an index entry for a new row, but defer any uniqueness-violation error until end of statement or even later. To avoid unnecessary repeat searches of the index, the index access method should do a preliminary uniqueness check during the initial insertion. If this shows that there is definitely no conflicting live
tupple, we are done. Otherwise, we schedule a recheck to occur when it is time to enforce the constraint. If, at the time of the recheck, both the inserted tuple and some other tuple with the same key are live, then the error must be reported. (Note that for this purpose, “live” actually means “any tuple in the index entry’s HOT chain is live”.) To implement this, the aminsert function is passed a checkUnique parameter having one of the following values:

- **UNIQUE_CHECK_NO** indicates that no uniqueness checking should be done (this is not a unique index).

- **UNIQUE_CHECK_YES** indicates that this is a non-deferrable unique index, and the uniqueness check must be done immediately, as described above.

- **UNIQUE_CHECK_PARTIAL** indicates that the unique constraint is deferrable. PostgreSQL will use this mode to insert each row’s index entry. The access method must allow duplicate entries into the index, and report any potential duplicates by returning false from aminsert. For each row for which false is returned, a deferred recheck will be scheduled.

The access method must identify any rows which might violate the unique constraint, but it is not an error for it to report false positives. This allows the check to be done without waiting for other transactions to finish; conflicts reported here are not treated as errors and will be rechecked later, by which time they may no longer be conflicts.

- **UNIQUE_CHECK_EXISTING** indicates that this is a deferred recheck of a row that was reported as a potential uniqueness violation. Although this is implemented by calling aminsert, the access method must not insert a new index entry in this case. The index entry is already present. Rather, the access method must check to see if there is another live index entry. If so, and if the target row is also still live, report error.

It is recommended that in a **UNIQUE_CHECK_EXISTING** call, the access method further verify that the target row actually does have an existing entry in the index, and report error if not. This is a good idea because the index tuple values passed to aminsert will have been recomputed. If the index definition involves functions that are not really immutable, we might be checking the wrong area of the index. Checking that the target row is found in the recheck verifies that we are scanning for the same tuple values as were used in the original insertion.

### 61.6. Index Cost Estimation Functions

The amcostestimate function is given information describing a possible index scan, including lists of WHERE and ORDER BY clauses that have been determined to be usable with the index. It must return estimates of the cost of accessing the index and the selectivity of the WHERE clauses (that is, the fraction of parent-table rows that will be retrieved during the index scan). For simple cases, nearly all the work of the cost estimator can be done by calling standard routines in the optimizer; the point of having an amcostestimate function is to allow index access methods to provide index-type-specific knowledge, in case it is possible to improve on the standard estimates.

Each amcostestimate function must have the signature:

```c
void amcostestimate (PlannerInfo *root,
    IndexPath *path,
    double loop_count,
    Cost *indexStartupCost,
    Cost *indexTotalCost,
    Selectivity *indexSelectivity,
    double *indexCorrelation,
    double *indexPages);
```

The first three parameters are inputs:
**Index Access Method**  
**Interface Definition**  

*root*

The planner's information about the query being processed.

*path*

The index access path being considered. All fields except cost and selectivity values are valid.

*loop_count*

The number of repetitions of the index scan that should be factored into the cost estimates. This will typically be greater than one when considering a parameterized scan for use in the inside of a nestloop join. Note that the cost estimates should still be for just one scan; a larger *loop_count* means that it may be appropriate to allow for some caching effects across multiple scans.

The last five parameters are pass-by-reference outputs:

*indexStartupCost*

Set to cost of index start-up processing

*indexTotalCost*

Set to total cost of index processing

*indexSelectivity*

Set to index selectivity

*indexCorrelation*

Set to correlation coefficient between index scan order and underlying table's order

*indexPages*

Set to number of index leaf pages

Note that cost estimate functions must be written in C, not in SQL or any available procedural language, because they must access internal data structures of the planner/optimizer.

The index access costs should be computed using the parameters used by `src/backend/optimizer/path/costsize.c`: a sequential disk block fetch has cost `seq_page_cost`, a non-sequential fetch has cost `random_page_cost`, and the cost of processing one index row should usually be taken as `cpu_index_tuple_cost`. In addition, an appropriate multiple of `cpu_operator_cost` should be charged for any comparison operators invoked during index processing (especially evaluation of the indexquals themselves).

The access costs should include all disk and CPU costs associated with scanning the index itself, but *not* the costs of retrieving or processing the parent-table rows that are identified by the index.

The “start-up cost” is the part of the total scan cost that must be expended before we can begin to fetch the first row. For most indexes this can be taken as zero, but an index type with a high start-up cost might want to set it nonzero.

The *indexSelectivity* should be set to the estimated fraction of the parent table rows that will be retrieved during the index scan. In the case of a lossy query, this will typically be higher than the fraction of rows that actually pass the given qual conditions.

The *indexCorrelation* should be set to the correlation (ranging between -1.0 and 1.0) between the index order and the table order. This is used to adjust the estimate for the cost of fetching rows from the parent table.
The *indexPages* should be set to the number of leaf pages. This is used to estimate the number of workers for parallel index scan.

When *loop_count* is greater than one, the returned numbers should be averages expected for any one scan of the index.

**Cost Estimation**

A typical cost estimator will proceed as follows:

1. Estimate and return the fraction of parent-table rows that will be visited based on the given qual conditions. In the absence of any index-type-specific knowledge, use the standard optimizer function *clauselist_selectivity()*:

   ```c
   *indexSelectivity = clauselist_selectivity(root, path->indexquals, path->indexinfo->rel->relid, JOIN_INNER, NULL);
   ```

2. Estimate the number of index rows that will be visited during the scan. For many index types this is the same as *indexSelectivity* times the number of rows in the index, but it might be more. (Note that the index's size in pages and rows is available from the *path->indexinfo* struct.)

3. Estimate the number of index pages that will be retrieved during the scan. This might be just *indexSelectivity* times the index's size in pages.

4. Compute the index access cost. A generic estimator might do this:

   ```c
   /*
   * Our generic assumption is that the index pages will be read
   * sequentially, so they cost seq_page_cost each, not
   * random_page_cost.
   * Also, we charge for evaluation of the indexquals at each
   * index row.
   * All the costs are assumed to be paid incrementally during
   * the scan.
   */
   cost_qual_eval(&index_qual_cost, path->indexquals, root);
   *indexStartupCost = index_qual_cost.startup;
   *indexTotalCost = seq_page_cost * numIndexPages +
                   (cpu_index_tuple_cost + index_qual_cost.per_tuple) *
                   numIndexTuples;
   ```

   However, the above does not account for amortization of index reads across repeated index scans.

5. Estimate the index correlation. For a simple ordered index on a single field, this can be retrieved from *pg_statistic*. If the correlation is not known, the conservative estimate is zero (no correlation).

Examples of cost estimator functions can be found in *src/backend/utils/adt/selfuncs.c*. 

2201
Chapter 62. Generic WAL Records

Although all built-in WAL-logged modules have their own types of WAL records, there is also a generic WAL record type, which describes changes to pages in a generic way. This is useful for extensions that provide custom access methods, because they cannot register their own WAL redo routines.

The API for constructing generic WAL records is defined in access/generic_xlog.h and implemented in access/transam/generic_xlog.c.

To perform a WAL-logged data update using the generic WAL record facility, follow these steps:

1. `state = GenericXLogStart(relation)` — start construction of a generic WAL record for the given relation.

2. `page = GenericXLogRegisterBuffer(state, buffer, flags)` — register a buffer to be modified within the current generic WAL record. This function returns a pointer to a temporary copy of the buffer's page, where modifications should be made. (Do not modify the buffer's contents directly.) The third argument is a bit mask of flags applicable to the operation. Currently the only such flag is `GENERIC_XLOG_FULL_IMAGE`, which indicates that a full-page image rather than a delta update should be included in the WAL record. Typically this flag would be set if the page is new or has been rewritten completely. `GenericXLogRegisterBuffer` can be repeated if the WAL-logged action needs to modify multiple pages.

3. Apply modifications to the page images obtained in the previous step.

4. `GenericXLogFinish(state)` — apply the changes to the buffers and emit the generic WAL record.

WAL record construction can be canceled between any of the above steps by calling `GenericXLogAbort(state)`. This will discard all changes to the page image copies.

Please note the following points when using the generic WAL record facility:

- No direct modifications of buffers are allowed! All modifications must be done in copies acquired from `GenericXLogRegisterBuffer()`. In other words, code that makes generic WAL records should never call `BufferGetPage()` for itself. However, it remains the caller's responsibility to pin/unpin and lock/unlock the buffers at appropriate times. Exclusive lock must be held on each target buffer from before `GenericXLogRegisterBuffer()` until after `GenericXLogFinish()`.

- Registrations of buffers (step 2) and modifications of page images (step 3) can be mixed freely, i.e., both steps may be repeated in any sequence. Keep in mind that buffers should be registered in the same order in which locks are to be obtained on them during replay.

- The maximum number of buffers that can be registered for a generic WAL record is `MAX_GENERIC_XLOG_PAGES`. An error will be thrown if this limit is exceeded.

- Generic WAL assumes that the pages to be modified have standard layout, and in particular that there is no useful data between `pd_lower` and `pd_upper`.

- Since you are modifying copies of buffer pages, `GenericXLogStart()` does not start a critical section. Thus, you can safely do memory allocation, error throwing, etc. between `GenericXLogStart()` and `GenericXLogFinish()`. The only actual critical section is present inside `GenericXLogFinish()`. There is no need to worry about calling `GenericXLogAbort()` during an error exit, either.

- `GenericXLogFinish()` takes care of marking buffers dirty and setting their LSNs. You do not need to do this explicitly.
Generic WAL Records

- For unlogged relations, everything works the same except that no actual WAL record is emitted. Thus, you typically do not need to do any explicit checks for unlogged relations.

- The generic WAL redo function will acquire exclusive locks to buffers in the same order as they were registered. After redoing all changes, the locks will be released in the same order.

- If GENERIC_XLOG_FULL_IMAGE is not specified for a registered buffer, the generic WAL record contains a delta between the old and the new page images. This delta is based on byte-by-byte comparison. This is not very compact for the case of moving data within a page, and might be improved in the future.
Chapter 63. B-Tree Indexes

63.1. Introduction

PostgreSQL includes an implementation of the standard btree (multi-way binary tree) index data structure. Any data type that can be sorted into a well-defined linear order can be indexed by a btree index. The only limitation is that an index entry cannot exceed approximately one-third of a page (after TOAST compression, if applicable).

Because each btree operator class imposes a sort order on its data type, btree operator classes (or, really, operator families) have come to be used as PostgreSQL's general representation and understanding of sorting semantics. Therefore, they've acquired some features that go beyond what would be needed just to support btree indexes, and parts of the system that are quite distant from the btree AM make use of them.

63.2. Behavior of B-Tree Operator Classes

As shown in Table 38.2, a btree operator class must provide five comparison operators, \(<\), \(\leq\), \(=\), \(\geq\), and \(>\). One might expect that \(\neq\) should also be part of the operator class, but it is not, because it would almost never be useful to use a \(\neq\) WHERE clause in an index search. (For some purposes, the planner treats \(\neq\) as associated with a btree operator class; but it finds that operator via the \(=\) operator's negator link, rather than from \texttt{pg_amop}.)

When several data types share near-identical sorting semantics, their operator classes can be grouped into an operator family. Doing so is advantageous because it allows the planner to make deductions about cross-type comparisons. Each operator class within the family should contain the single-type operators (and associated support functions) for its input data type, while cross-type comparison operators and support functions are “loose” in the family. It is recommendable that a complete set of cross-type operators be included in the family, thus ensuring that the planner can represent any comparison conditions that it deduces from transitivity.

There are some basic assumptions that a btree operator family must satisfy:

- An \(=\) operator must be an equivalence relation; that is, for all non-null values \(A\), \(B\), \(C\) of the data type:
  
  - \(A = A\) is true (reflexive law)
  - if \(A = B\), then \(B = A\) (symmetric law)
  - if \(A = B\) and \(B = C\), then \(A = C\) (transitive law)

- A \(<\) operator must be a strong ordering relation; that is, for all non-null values \(A\), \(B\), \(C\):
  
  - \(A < A\) is false (irreflexive law)
  - if \(A < B\) and \(B < C\), then \(A < C\) (transitive law)
  - Furthermore, the ordering is total; that is, for all non-null values \(A\), \(B\):
    
    - exactly one of \(A < B\), \(A = B\), and \(B < A\) is true (trichotomy law)

(The trichotomy law justifies the definition of the comparison support function, of course.)

The other three operators are defined in terms of \(=\) and \(<\) in the obvious way, and must act consistently with them.

For an operator family supporting multiple data types, the above laws must hold when \(A\), \(B\), \(C\) are taken from any data types in the family. The transitive laws are the trickiest to ensure, as in cross-type situations they represent statements that the behaviors of two or three different operators are consistent.
As an example, it would not work to put float8 and numeric into the same operator family, at least not with the current semantics that numeric values are converted to float8 for comparison to a float8. Because of the limited accuracy of float8, this means there are distinct numeric values that will compare equal to the same float8 value, and thus the transitive law would fail.

Another requirement for a multiple-data-type family is that any implicit or binary-coercion casts that are defined between data types included in the operator family must not change the associated sort ordering.

It should be fairly clear why a btree index requires these laws to hold within a single data type: without them there is no ordering to arrange the keys with. Also, index searches using a comparison key of a different data type require comparisons to behave sanely across two data types. The extensions to three or more data types within a family are not strictly required by the btree index mechanism itself, but the planner relies on them for optimization purposes.

### 63.3. B-Tree Support Functions

As shown in Table 38.8, btree defines one required and two optional support functions.

For each combination of data types that a btree operator family provides comparison operators for, it must provide a comparison support function, registered in pg_amproc with support function number 1 and amproclefttype/amprocrighttype equal to the left and right data types for the comparison (i.e., the same data types that the matching operators are registered with in pg_amop). The comparison function must take two non-null values \( A \) and \( B \) and return an int32 value that is \(< 0, 0, \) or \(> 0 \) when \( A < B, A = B, \) or \( A > B \), respectively. A null result is disallowed: all values of the data type must be comparable. See src/backend/access/nbtree/nbtcompare.c for examples.

If the compared values are of a collatable data type, the appropriate collation OID will be passed to the comparison support function, using the standard PG_GET_COLLATION() mechanism.

Optionally, a btree operator family may provide sort support function(s), registered under support function number 2. These functions allow implementing comparisons for sorting purposes in a more efficient way than naively calling the comparison support function. The APIs involved in this are defined in src/include/utils/sortsupport.h.

Optionally, a btree operator family may provide in_range support function(s), registered under support function number 3. These are not used during btree index operations; rather, they extend the semantics of the operator family so that it can support window clauses containing the RANGE PRECEDING and RANGE offset FOLLOWING frame bound types (see Section 4.2.8). Fundamentally, the extra information provided is how to add or subtract an offset value in a way that is compatible with the family's data ordering.

An in_range function must have the signature

\[
\text{in\_range}(\text{val } \text{type1}, \text{ base } \text{type1}, \text{ offset } \text{type2}, \text{ sub } \text{bool}, \text{ less } \text{bool})
\]

returns bool

\( \text{val} \) and \( \text{base} \) must be of the same type, which is one of the types supported by the operator family (i.e., a type for which it provides an ordering). However, \( \text{offset} \) could be of a different type, which might be one otherwise unsupported by the family. An example is that the built-in time_ops family provides an in_range function that has \( \text{offset} \) of type interval. A family can provide in_range functions for any of its supported types and one or more \( \text{offset} \) types. Each in_range function should be entered in pg_amproc with amproclefttype equal to \( \text{type1} \) and amprocrighttype equal to \( \text{type2} \).

The essential semantics of an in_range function depend on the two boolean flag parameters. It should add or subtract \( \text{base} \) and \( \text{offset} \), then compare \( \text{val} \) to the result, as follows:
• if !sub and !less, return val >= (base + offset)
• if !sub and less, return val <= (base + offset)
• if sub and !less, return val >= (base - offset)
• if sub and less, return val <= (base - offset)

Before doing so, the function should check the sign of offset: if it is less than zero, raise error
ERRCODE_INVALID_PRECEDING_OR_FOLLOWING_SIZE (22013) with error text like “invalid
preceding or following size in window function”. (This is required by the SQL standard, although
nonstandard operator families might perhaps choose to ignore this restriction, since there seems to be
little semantic necessity for it.) This requirement is delegated to the in_range function so that the
core code needn’t understand what “less than zero” means for a particular data type.

An additional expectation is that in_range functions should, if practical, avoid throwing an error
if base + offset or base - offset would overflow. The correct comparison result can be de-
termined even if that value would be out of the data type’s range. Note that if the data type includes
concepts such as “infinity” or “NaN”, extra care may be needed to ensure that in_range’s results
agree with the normal sort order of the operator family.

The results of the in_range function must be consistent with the sort ordering imposed by the
operator family. To be precise, given any fixed values of offset and sub, then:
• If in_range with less = true is true for some val1 and base, it must be true for every val2
  <= val1 with the same base.
• If in_range with less = true is false for some val1 and base, it must be false for every val2
  >= val1 with the same base.
• If in_range with less = true is true for some val and base1, it must be true for every base2
  >= base1 with the same val.
• If in_range with less = true is false for some val and base1, it must be false for every base2
  <= base1 with the same val.

Analogous statements with inverted conditions hold when less = false.

If the type being ordered (type1) is collatable, the appropriate collation OID will be passed to the
in_range function, using the standard PG_GET_COLLATION() mechanism.

in_range functions need not handle NULL inputs, and typically will be marked strict.

### 63.4. Implementation

An introduction to the btree index implementation can be found in src/backend/access/nbtree/README.
Chapter 64. GiST Indexes

64.1. Introduction

GiST stands for Generalized Search Tree. It is a balanced, tree-structured access method, that acts as a base template in which to implement arbitrary indexing schemes. B-trees, R-trees and many other indexing schemes can be implemented in GiST.

One advantage of GiST is that it allows the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert.

Some of the information here is derived from the University of California at Berkeley's GiST Indexing Project web site¹ and Marcel Kornacker's thesis, Access Methods for Next-Generation Database Systems². The GiST implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov, and there is more information on their web site³.

64.2. Built-in Operator Classes

The core PostgreSQL distribution includes the GiST operator classes shown in Table 64.1. (Some of the optional modules described in Appendix F provide additional GiST operator classes.)

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
<th>Ordering Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>box_ops</td>
<td>box</td>
<td>&amp;&amp;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&amp;&gt;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td>circle_ops</td>
<td>circle</td>
<td>&amp;&amp;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&amp;&gt;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td>inet_ops</td>
<td>inet, cidr</td>
<td>&amp;&amp;</td>
<td>&gt;&gt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&lt;=</td>
</tr>
<tr>
<td>point_ops</td>
<td>point</td>
<td>&gt;&gt;</td>
<td>&gt;&gt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>&lt;=</td>
</tr>
<tr>
<td>poly_ops</td>
<td>polygon</td>
<td>&amp;&amp;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&amp;&gt;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td>range_ops</td>
<td>any range type</td>
<td>&amp;&amp;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&amp;&gt;</td>
<td>&amp;&lt;</td>
</tr>
<tr>
<td>tsquery_ops</td>
<td>tsquery</td>
<td>&lt;&lt;</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>tsvector_ops</td>
<td>tsvector</td>
<td>@</td>
<td></td>
</tr>
</tbody>
</table>

For historical reasons, the inet_ops operator class is not the default class for types inet and cidr. To use it, mention the class name in CREATE INDEX, for example

```sql
CREATE INDEX ON my_table USING GIST (my_inet_column inet_ops);
```

¹ http://gist.cs.berkeley.edu/
² http://www.sai.msu.su/~megera/postgres/gist/papers/concurrency/access-methods-for-next-generation.pdf.gz
³ http://www.sai.msu.su/~megera/postgres/gist/
64.3. Extensibility

Traditionally, implementing a new index access method meant a lot of difficult work. It was necessary to understand the inner workings of the database, such as the lock manager and Write-Ahead Log. The GiST interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The GiST layer itself takes care of concurrency, logging and searching the tree structure.

This extensibility should not be confused with the extensibility of the other standard search trees in terms of the data they can handle. For example, PostgreSQL supports extensible B-trees and hash indexes. That means that you can use PostgreSQL to build a B-tree or hash over any data type you want. But B-trees only support range predicates (\(<\), \(\leq\), \(\geq\)), and hash indexes only support equality queries.

So if you index, say, an image collection with a PostgreSQL B-tree, you can only issue queries such as “is imagex equal to imagey”, “is imagex less than imagey” and “is imagex greater than imagey”. Depending on how you define “equals”, “less than” and “greater than” in this context, this could be useful. However, by using a GiST based index, you could create ways to ask domain-specific questions, perhaps “find all images of horses” or “find all over-exposed images”.

All it takes to get a GiST access method up and running is to implement several user-defined methods, which define the behavior of keys in the tree. Of course these methods have to be pretty fancy to support fancy queries, but for all the standard queries (B-trees, R-trees, etc.) they're relatively straightforward. In short, GiST combines extensibility along with generality, code reuse, and a clean interface.

There are five methods that an index operator class for GiST must provide, and four that are optional. Correctness of the index is ensured by proper implementation of the \texttt{same}, \texttt{consistent} and \texttt{union} methods, while efficiency (size and speed) of the index will depend on the \texttt{penalty} and \texttt{picksplit} methods. Two optional methods are \texttt{compress} and \texttt{decompress}, which allow an index to have internal tree data of a different type than the data it indexes. The leaves are to be of the indexed data type, while the other tree nodes can be of any C struct (but you still have to follow PostgreSQL data type rules here, see about \texttt{varlena} for variable sized data). If the tree's internal data type exists at the SQL level, the \texttt{STORAGE} option of the \texttt{CREATE OPERATOR CLASS} command can be used. The optional eighth method is \texttt{distance}, which is needed if the operator class wishes to support ordered scans (nearest-neighbor searches). The optional ninth method \texttt{fetch} is needed if the operator class wishes to support index-only scans, except when the \texttt{compress} method is omitted.

Given an index entry \(p\) and a query value \(q\), this function determines whether the index entry is "consistent" with the query; that is, could the predicate "\texttt{indexed_column indexable_operator q}" be true for any row represented by the index entry? For a leaf index entry this is equivalent to testing the indexable condition, while for an internal tree node this determines whether it is necessary to scan the subtree of the index represented by the tree node. When the result is true, a recheck flag must also be returned. This indicates whether the predicate is certainly true or only possibly true. If \texttt{recheck = false} then the index has tested the predicate condition exactly, whereas if \texttt{recheck = true} the row is only a candidate match. In that case the system will automatically evaluate the \texttt{indexable_operator} against the actual row value to see if it is really a match. This convention allows GiST to support both lossless and lossy index structures.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_consistent(internal, data_type, smallint, oid, internal) RETURNS bool AS 'MODULE_PATHNAME' LANGUAGE C STRICT;
```

And the matching code in the C module could then follow this skeleton:
GiST Indexes

Datum
my_consistent(PG_FUNCTION_ARGS)
{
    GISTENTRY   *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
data_type    *query = PG_GETARG_DATA_TYPE_P(1);
StrategyNumber strategy = (StrategyNumber)
PG_GETARG_UINT16(2);
    /* Oid subtype = PG_GETARG_OID(3); */
    bool       *recheck = (bool *) PG_GETARG_POINTER(4);
data_type    *key = DatumGetDataType(entry->key);
bool          retval;

    /*
     * determine return value as a function of strategy, key and
     * query.
     *
     * Use GIST_LEAF(entry) to know where you're called in the
     * index tree,
     * which comes handy when supporting the = operator for
     * example (you could
     * check for non empty union() in non-leaf nodes and
     * equality in leaf
     * nodes).
     */

    *recheck = true;        /* or false if check is exact */

    PG_RETURN_BOOL(retval);
}

Here, key is an element in the index and query the value being looked up in the index. The StrategyNumber parameter indicates which operator of your operator class is being applied — it matches one of the operator numbers in the CREATE OPERATOR CLASS command.

Depending on which operators you have included in the class, the data type of query could vary with the operator, since it will be whatever type is on the righthand side of the operator, which might be different from the indexed data type appearing on the lefthand side. (The above code skeleton assumes that only one type is possible; if not, fetching the query argument value would have to depend on the operator.) It is recommended that the SQL declaration of the consistent function use the opclass's indexed data type for the query argument, even though the actual type might be something else depending on the operator.

union

This method consolidates information in the tree. Given a set of entries, this function generates a new index entry that represents all the given entries.

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_union(internal, internal)
RETURNS storage_type
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:
GiST Indexes

```
PG_FUNCTION_INFO_V1(my_union);

Datum
my_union(PG_FUNCTION_ARGS)
{
    GistEntryVector *entryvec = (GistEntryVector *)
    PG_GETARG_POINTER(0);
    GISTENTRY *ent = entryvec->vector;
    data_type *out,
        *tmp,
        *old;
    int         numranges,
                i = 0;

    numranges = entryvec->n;
    tmp = DatumGetDataType(ent[0].key);
    out = tmp;

    if (numranges == 1)
        { out = data_type_deep_copy(tmp);
          PG_RETURN_DATA_TYPE_P(out);
        }
    for (i = 1; i < numranges; i++)
        { old = out;
          tmp = DatumGetDataType(ent[i].key);
          out = my_union_implementation(out, tmp);
        }
    PG_RETURN_DATA_TYPE_P(out);
}
```

As you can see, in this skeleton we're dealing with a data type where \texttt{union(X, Y, Z) = union(union(X, Y), Z)}. It's easy enough to support data types where this is not the case, by implementing the proper union algorithm in this GiST support method.

The result of the \texttt{union} function must be a value of the index's storage type, whatever that is (it might or might not be different from the indexed column's type). The \texttt{union} function should return a pointer to newly \texttt{palloc}()ed memory. You can't just return the input value as-is, even if there is no type change.

As shown above, the \texttt{union} function's first \texttt{internal} argument is actually a GistEntryVector pointer. The second argument is a pointer to an integer variable, which can be ignored. (It used to be required that the \texttt{union} function store the size of its result value into that variable, but this is no longer necessary.)

\textbf{compress}

Converts a data item into a format suitable for physical storage in an index page. If the \texttt{compress} method is omitted, data items are stored in the index without modification.

The SQL declaration of the function must look like this:

```
CREATE OR REPLACE FUNCTION my_compress(internal)
```
GiST Indexes

RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_compress);

Datum
my_compress(PG_FUNCTION_ARGS)
{
    GISTENTRY *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
    GISTENTRY *retval;

    if (entry->leafkey)
    {
        /* replace entry->key with a compressed version */
        compressed_data_type *compressed_data = palloc(sizeof(compressed_data_type));

        /* fill *compressed_data from entry->key ... */
        retval = palloc(sizeof(GISTENTRY));
        gistentryinit(*retval, PointerGetDatum(compressed_data),
                      entry->rel, entry->page, entry->offset,
                      FALSE);
    }
    else
    {
        /* typically we needn't do anything with non-leaf
         * entries */
        retval = entry;
    }

    PG_RETURN_POINTER(retval);
}

You have to adapt compressed_data_type to the specific type you're converting to in order to compress your leaf nodes, of course.

decompress

Converts the stored representation of a data item into a format that can be manipulated by the other GiST methods in the operator class. If the decompress method is omitted, it is assumed that the other GiST methods can work directly on the stored data format. (decompress is not necessarily the reverse of the compress method; in particular, if compress is lossy then it's impossible for decompress to exactly reconstruct the original data. decompress is not necessarily equivalent to fetch, either, since the other GiST methods might not require full reconstruction of the data.)

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_decompress(internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:
GiST Indexes

PG_FUNCTION_INFO_V1(my_decompress);

Datum
my_decompress(PG_FUNCTION_ARGS)
{
    PG_RETURN_POINTER(PG_GETARG_POINTER(0));
}

The above skeleton is suitable for the case where no decompression is needed. (But, of course, omitting the method altogether is even easier, and is recommended in such cases.)

penalty

Returns a value indicating the “cost” of inserting the new entry into a particular branch of the tree. Items will be inserted down the path of least penalty in the tree. Values returned by penalty should be non-negative. If a negative value is returned, it will be treated as zero.

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_penalty(internal, internal, internal)
    RETURNS internal
    AS 'MODULE_PATHNAME'
    LANGUAGE C STRICT;  -- in some cases penalty functions need not be strict

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_penalty);

Datum
my_penalty(PG_FUNCTION_ARGS)
{
    GISTENTRY  *origentry = (GISTENTRY *) PG_GETARG_POINTER(0);
    GISTENTRY  *newentry = (GISTENTRY *) PG_GETARG_POINTER(1);
    float      *penalty = (float *) PG_GETARG_POINTER(2);
    data_type  *orig = DatumGetDataType(origentry->key);
    data_type  *new = DatumGetDataType(newentry->key);

    *penalty = my_penalty_implementation(orig, new);
    PG_RETURN_POINTER(penalty);
}

For historical reasons, the penalty function doesn't just return a float result; instead it has to store the value at the location indicated by the third argument. The return value per se is ignored, though it's conventional to pass back the address of that argument.

The penalty function is crucial to good performance of the index. It'll get used at insertion time to determine which branch to follow when choosing where to add the new entry in the tree. At query time, the more balanced the index, the quicker the lookup.

picksplit

When an index page split is necessary, this function decides which entries on the page are to stay on the old page, and which are to move to the new page.

The SQL declaration of the function must look like this:
CREATE OR REPLACE FUNCTION my_picksplit(internal, internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_picksplit);

Datum
my_picksplit(PG_FUNCTION_ARGS)
{
    GistEntryVector *entryvec = (GistEntryVector *)
    PG_GETARG_POINTER(0);
    GIST_SPLITVEC *v = (GIST_SPLITVEC *) PG_GETARG_POINTER(1);
    OffsetNumber maxoff = entryvec->n - 1;
    GISTENTRY *ent = entryvec->vector;
    int i,
        nbytes;
    OffsetNumber *left,
        *right;
    data_type *tmp_union;
    data_type *unionL;
    data_type *unionR;
    GISTENTRY **raw_entryvec;

    maxoff = entryvec->n - 1;
    nbytes = (maxoff + 1) * sizeof(OffsetNumber);
    v->spl_left = (OffsetNumber *) palloc(nbytes);
    left = v->spl_left;
    v->spl_nleft = 0;
    v->spl_right = (OffsetNumber *) palloc(nbytes);
    right = v->spl_right;
    v->spl_nright = 0;
    unionL = NULL;
    unionR = NULL;

    /* Initialize the raw entry vector. */
    raw_entryvec = (GISTENTRY **) malloc(entryvec->n *
        sizeof(void *));
    for (i = FirstOffsetNumber; i <= maxoff; i =
        OffsetNumberNext(i))
        raw_entryvec[i] = &(entryvec->vector[i]);

    for (i = FirstOffsetNumber; i <= maxoff; i =
        OffsetNumberNext(i))
        { int real_index = raw_entryvec[i] - entryvec->vector;
            tmp_union = DatumGetDataType(entryvec->vector[real_index].key);
            Assert(tmp_union != NULL);
/*
 * Choose where to put the index entries and update unionL and unionR
 * accordingly. Append the entries to either v_spl_left or
 * v_spl_right, and care about the counters.
 */

if (my_choice_is_left(unionL, curl, unionR, curr))
{
    if (unionL == NULL)
        unionL = tmp_union;
    else
        unionL = my_union_implementation(unionL,
            tmp_union);

    *(v->spl_nleft);
} else
{
    /*
     * Same on the right
     */
}

v->spl_ldatum = DataTypeGetDatum(unionL);
v->spl_rdatum = DataTypeGetDatum(unionR);
PG_RETURN_POINTER(v);

Notice that the picksplit function’s result is delivered by modifying the passed-in v structure.
The return value per se is ignored, though it’s conventional to pass back the address of v.

Like penalty, the picksplit function is crucial to good performance of the index. Designing
suitable penalty and picksplit implementations is where the challenge of implementing
well-performing GiST indexes lies.

same

Returns true if two index entries are identical, false otherwise. (An “index entry” is a value of the
index’s storage type, not necessarily the original indexed column’s type.)

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_same(storage_type, storage_type,
    internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:
Datum
my_same(PG_FUNCTION_ARGS)
{
    prefix_range *v1 = PG_GETARG_PREFIX_RANGE_P(0);
    prefix_range *v2 = PG_GETARG_PREFIX_RANGE_P(1);
    bool *result = (bool *) PG_GETARG_POINTER(2);

    *result = my_eq(v1, v2);
    PG_RETURN_POINTER(result);
}

For historical reasons, the `same` function doesn’t just return a Boolean result; instead it has to store the flag at the location indicated by the third argument. The return value per se is ignored, though it’s conventional to pass back the address of that argument.

**distance**

Given an index entry $p$ and a query value $q$, this function determines the index entry’s “distance” from the query value. This function must be supplied if the operator class contains any ordering operators. A query using the ordering operator will be implemented by returning index entries with the smallest “distance” values first, so the results must be consistent with the operator’s semantics. For a leaf index entry the result just represents the distance to the index entry; for an internal tree node, the result must be the smallest distance that any child entry could have.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_distance(internal, data_type, smallint, oid, internal)
RETURNS float8
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

And the matching code in the C module could then follow this skeleton:

```c
PG_FUNCTION_INFO_V1(my_distance);

Datum
my_distance(PG_FUNCTION_ARGS)
{
    GISTENTRY *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
    data_type *query = PG_GETARG_DATA_TYPE_P(1);
    StrategyNumber strategy = (StrategyNumber)
    PG_GETARG_UINT16(2);
    /* Oid subtype = PG_GETARG_OID(3); */
    /* bool *recheck = (bool *) PG_GETARG_POINTER(4); */
    data_type *key = DatumGetDataType(entry->key);
    double retval;

    /*
     * determine return value as a function of strategy, key and query.
     */
    PG_RETURN_FLOAT8(retval);
}
```

The arguments to the `distance` function are identical to the arguments of the `consistent` function.
Some approximation is allowed when determining the distance, so long as the result is never
greater than the entry's actual distance. Thus, for example, distance to a bounding box is usually
sufficient in geometric applications. For an internal tree node, the distance returned must not
be greater than the distance to any of the child nodes. If the returned distance is not exact, the
function must set \*recheck to true. (This is not necessary for internal tree nodes; for them, the
calculation is always assumed to be inexact.) In this case the executor will calculate the accurate
distance after fetching the tuple from the heap, and reorder the tuples if necessary.

If the distance function returns \*recheck = true for any leaf node, the original ordering
operator's return type must be float8 or float4, and the distance function's result values
must be comparable to those of the original ordering operator, since the executor will sort using
both distance function results and recalculated ordering-operator results. Otherwise, the distance
function's result values can be any finite float8 values, so long as the relative order of the
result values matches the order returned by the ordering operator. (Infinity and minus infinity are
used internally to handle cases such as nulls, so it is not recommended that distance functions
return these values.)

**fetch**

Converts the compressed index representation of a data item into the original data type, for in-
dex-only scans. The returned data must be an exact, non-lossy copy of the originally indexed
value.

The SQL declaration of the function must look like this:

```
CREATE OR REPLACE FUNCTION my_fetch(internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

The argument is a pointer to a GISTENTRY struct. On entry, its key field contains a non-NULL
leaf datum in compressed form. The return value is another GISTENTRY struct, whose key field
contains the same datum in its original, uncompressed form. If the opclass's compress function
does nothing for leaf entries, the fetch method can return the argument as-is. Or, if the opclass
does not have a compress function, the fetch method can be omitted as well, since it would
necessarily be a no-op.

The matching code in the C module could then follow this skeleton:

```
PG_FUNCTION_INFO_V1(my_fetch);

Datum
my_fetch(PG_FUNCTION_ARGS)
{
    GISTENTRY  *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
    input_data_type *in = DatumGetP(entry->key);
    fetched_data_type *fetched_data;
    GISTENTRY  *retval;

    retval = palloc(sizeof(GISTENTRY));
    fetched_data = palloc(sizeof(fetched_data_type));

    /*
     * Convert 'fetched_data' into the a Datum of the original
datatype.
     */
```
/* fill *retval from fetch_data. */
gistentryinit(*retval, PointerGetDatum(converted_datum),
    entry->rel, entry->page, entry->offset,
    FALSE);

PG_RETURN_POINTER(retval);
}

If the compress method is lossy for leaf entries, the operator class cannot support index-only scans,
and must not define a fetch function.

All the GiST support methods are normally called in short-lived memory contexts; that is, Current-
MemoryContext will get reset after each tuple is processed. It is therefore not very important to
worry about pfree'ing everything you palloc. However, in some cases it's useful for a support method
to cache data across repeated calls. To do that, allocate the longer-lived data in fcinfo->flinfo-
>fn_mcxt, and keep a pointer to it in fcinfo->flinfo->fn_extra. Such data will survive for
the life of the index operation (e.g., a single GiST index scan, index build, or index tuple insertion).
Be careful to pfree the previous value when replacing a fn_extra value, or the leak will accumulate
for the duration of the operation.

64.4. Implementation

64.4.1. GiST buffering build

Building large GiST indexes by simply inserting all the tuples tends to be slow, because if the index
tuples are scattered across the index and the index is large enough to not fit in cache, the insertions need
to perform a lot of random I/O. Beginning in version 9.2, PostgreSQL supports a more efficient method
to build GiST indexes based on buffering, which can dramatically reduce the number of random I/
Os needed for non-ordered data sets. For well-ordered data sets the benefit is smaller or non-existent,
because only a small number of pages receive new tuples at a time, and those pages fit in cache even
if the index as whole does not.

However, buffering index build needs to call the penalty function more often, which consumes
some extra CPU resources. Also, the buffers used in the buffering build need temporary disk space,
up to the size of the resulting index. Buffering can also influence the quality of the resulting index, in
both positive and negative directions. That influence depends on various factors, like the distribution
of the input data and the operator class implementation.

By default, a GiST index build switches to the buffering method when the index size reaches effec-
tive_cache_size. It can be manually turned on or off by the buffering parameter to the CREATE
INDEX command. The default behavior is good for most cases, but turning buffering off might speed
up the build somewhat if the input data is ordered.

64.5. Examples

The PostgreSQL source distribution includes several examples of index methods implemented us-
ing GiST. The core system currently provides text search support (indexing for tsvector and ts-
query) as well as R-Tree equivalent functionality for some of the built-in geometric data types (see
src/backend/access/gist/gistproc.c). The following contrib modules also contain
GiST operator classes:

btree_gist

B-tree equivalent functionality for several data types

cube

Indexing for multidimensional cubes
GiST Indexes

hstore
   Module for storing (key, value) pairs

intarray
   RD-Tree for one-dimensional array of int4 values

ltree
   Indexing for tree-like structures

pg_trgm
   Text similarity using trigram matching

seg
   Indexing for “float ranges”
Chapter 65. SP-GiST Indexes

65.1. Introduction

SP-GiST is an abbreviation for space-partitioned GiST. SP-GiST supports partitioned search trees, which facilitate development of a wide range of different non-balanced data structures, such as quad-trees, k-d trees, and radix trees (tries). The common feature of these structures is that they repeatedly divide the search space into partitions that need not be of equal size. Searches that are well matched to the partitioning rule can be very fast.

These popular data structures were originally developed for in-memory usage. In main memory, they are usually designed as a set of dynamically allocated nodes linked by pointers. This is not suitable for direct storing on disk, since these chains of pointers can be rather long which would require too many disk accesses. In contrast, disk-based data structures should have a high fanout to minimize I/O. The challenge addressed by SP-GiST is to map search tree nodes to disk pages in such a way that a search need access only a few disk pages, even if it traverses many nodes.

Like GiST, SP-GiST is meant to allow the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert.

Some of the information here is derived from Purdue University's SP-GiST Indexing Project website\(^1\). The SP-GiST implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov, and there is more information on their web site\(^2\).

65.2. Built-in Operator Classes

The core PostgreSQL distribution includes the SP-GiST operator classes shown in Table 65.1.

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>kd_point_ops</td>
<td>point</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>quad_point_ops</td>
<td>point</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>range_ops</td>
<td>any range type</td>
<td>&amp;&amp;</td>
</tr>
<tr>
<td>box_ops</td>
<td>box</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>poly_ops</td>
<td>polygon</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>text_ops</td>
<td>text</td>
<td>&lt;=</td>
</tr>
<tr>
<td>inet_ops</td>
<td>inet, cidr</td>
<td>&amp;&amp;</td>
</tr>
</tbody>
</table>

Of the two operator classes for type point, quad_point_ops is the default. kd_point_ops supports the same operators but uses a different index data structure which may offer better performance in some applications.

65.3. Extensibility

SP-GiST offers an interface with a high level of abstraction, requiring the access method developer to implement only methods specific to a given data type. The SP-GiST core is responsible for efficient

---

\(^1\) [https://www.cs.purdue.edu/spgist/](https://www.cs.purdue.edu/spgist/)

\(^2\) [http://www.sai.msu.su/~megera/wiki/spgist_dev](http://www.sai.msu.su/~megera/wiki/spgist_dev)
SP-GiST Indexes

disk mapping and searching the tree structure. It also takes care of concurrency and logging considerations.

Leaf tuples of an SP-GiST tree contain values of the same data type as the indexed column. Leaf tuples at the root level will always contain the original indexed data value, but leaf tuples at lower levels might contain only a compressed representation, such as a suffix. In that case the operator class support functions must be able to reconstruct the original value using information accumulated from the inner tuples that are passed through to reach the leaf level.

Inner tuples are more complex, since they are branching points in the search tree. Each inner tuple contains a set of one or more nodes, which represent groups of similar leaf values. A node contains a downlink that leads either to another, lower-level inner tuple, or to a short list of leaf tuples that all lie on the same index page. Each node normally has a label that describes it; for example, in a radix tree the node label could be the next character of the string value. (Alternatively, an operator class can omit the node labels, if it works with a fixed set of nodes for all inner tuples; see Section 65.4.2.) Optionally, an inner tuple can have a prefix value that describes all its members. In a radix tree this could be the common prefix of the represented strings. The prefix value is not necessarily really a prefix, but can be any data needed by the operator class; for example, in a quad-tree it can store the central point that the four quadrants are measured with respect to. A quad-tree inner tuple would then also contain four nodes corresponding to the quadrants around this central point.

Some tree algorithms require knowledge of level (or depth) of the current tuple, so the SP-GiST core provides the possibility for operator classes to manage level counting while descending the tree. There is also support for incrementally reconstructing the represented value when that is needed, and for passing down additional data (called traverse values) during a tree descent.

Note

The SP-GiST core code takes care of null entries. Although SP-GiST indexes do store entries for nulls in indexed columns, this is hidden from the index operator class code: no null index entries or search conditions will ever be passed to the operator class methods. (It is assumed that SP-GiST operators are strict and so cannot succeed for null values.) Null values are therefore not discussed further here.

There are five user-defined methods that an index operator class for SP-GiST must provide, and one is optional. All five mandatory methods follow the convention of accepting two internal arguments, the first of which is a pointer to a C struct containing input values for the support method, while the second argument is a pointer to a C struct where output values must be placed. Four of the mandatory methods just return void, since all their results appear in the output struct; but leaf_consistent additionally returns a boolean result. The methods must not modify any fields of their input structs. In all cases, the output struct is initialized to zeroes before calling the user-defined method. Optional sixth method compress accepts datum to be indexed as the only argument and returns value suitable for physical storage in leaf tuple.

The five mandatory user-defined methods are:

config

Returns static information about the index implementation, including the data type OIDs of the prefix and node label data types.

The SQL declaration of the function must look like this:

```
CREATE FUNCTION my_config(internal, internal) RETURNS void ... 
```

The first argument is a pointer to a spgConfigIn C struct, containing input data for the function. The second argument is a pointer to a spgConfigOut C struct, which the function must fill with result data.
typedef struct spgConfigIn
{
    Oid attType;       /* Data type to be indexed */
} spgConfigIn;

typedef struct spgConfigOut
{
    Oid prefixType;     /* Data type of inner-tuple prefixes */
    Oid labelType;      /* Data type of inner-tuple node labels */
    Oid leafType;       /* Data type of leaf-tuple values */
    bool canReturnData; /* Opclass can reconstruct original data */
    bool longValuesOK;  /* Opclass can cope with values > 1 page */
} spgConfigOut;

attType is passed in order to support polymorphic index operator classes; for ordinary fixed-data-type operator classes, it will always have the same value and so can be ignored.

For operator classes that do not use prefixes, prefixType can be set to VOIDOID. Likewise, for operator classes that do not use node labels, labelType can be set to VOIDOID. canReturnData should be set true if the operator class is capable of reconstructing the originally-supplied index value. longValuesOK should be set true only when the attType is of variable length and the operator class is capable of segmenting long values by repeated suffixing (see Section 65.4.1).

leafType is typically the same as attType. For the reasons of backward compatibility, method config can leave leafType uninitialized; that would give the same effect as setting leafType equal to attType. When attType and leafType are different, then optional method compress must be provided. Method compress is responsible for transformation of datums to be indexed from attType to leafType. Note: both consistent functions will get scankeys unchanged, without transformation using compress.

choose

Chooses a method for inserting a new value into an inner tuple.

The SQL declaration of the function must look like this:

CREATE FUNCTION my_choose(internal, internal) RETURNS void ...

The first argument is a pointer to a spgChooseIn C struct, containing input data for the function. The second argument is a pointer to a spgChooseOut C struct, which the function must fill with result data.

typedef struct spgChooseIn
{
    Datum datum;          /* original datum to be indexed */
    Datum leafDatum;      /* current datum to be stored at leaf */
    int level;            /* current level (counting from zero) */

    /* Data from current inner tuple */
bool allTheSame;    /* tuple is marked all-the-same? */
bool hasPrefix;     /* tuple has a prefix? */
Datum prefixDatum;  /* if so, the prefix value */
int nNodes;         /* number of nodes in the inner tuple */
Datum *nodeLabels;  /* node label values (NULL if none) */
} spgChooseIn;

typedef enum spgChooseResultType
{
    spgMatchNode = 1,             /* descend into existing node */
    spgAddNode,                   /* add a node to the inner tuple */
    spgSplitTuple                 /* split inner tuple (change its prefix) */
} spgChooseResultType;

typedef struct spgChooseOut
{
    spgChooseResultType resultType;    /* action code, see above */
    union
    {
        struct                  /* results for spgMatchNode */
        {
            int nodeN;        /* descend to this node (index from 0) */
            int levelAdd;    /* increment level by this much */
            Datum restDatum; /* new leaf datum */
        } matchNode;
    } struct
    {
        struct                  /* results for spgAddNode */
        {
            Datum nodeLabel; /* new node's label */
            int nodeN;       /* where to insert it (index from 0) */
        } addNode;
    } struct
    {
        /* Info to form new upper-level inner tuple with one child tuple */
        bool prefixHasPrefix; /* tuple should have a prefix? */
        Datum prefixPrefixDatum; /* if so, its value */
        int prefixNNodes;       /* number of nodes */
        Datum *prefixNodeLabels; /* their labels (or NULL for * no labels) */
        int childNodeN;         /* which node gets child tuple */
    } struct
    {
        /* Info to form new lower-level inner tuple with all old nodes */

bool postfixHasPrefix; /* tuple should have a prefix? */
Datum postfixPrefixDatum; /* if so, its value */
}
}
splitTuple;
}
result;
}
spgChooseOut;

datum is the original datum of spgConfigIn.attType type that was to be inserted into the index. leafDatum is a value of spgConfigOut.leafType type which is initially an result of method compress applied to datum when method compress is provided, or same value as datum otherwise. leafDatum can change at lower levels of the tree if the choose or picksplit methods change it. When the insertion search reaches a leaf page, the current value of leafDatum is what will be stored in the newly created leaf tuple. level is the current inner tuple's level, starting at zero for the root level. allTheSame is true if the current inner tuple is marked as containing multiple equivalent nodes (see Section 65.4.3). hasPrefix is true if the current inner tuple contains a prefix; if so, prefixDatum is its value. nNodes is the number of child nodes contained in the inner tuple, and nodeLabels is an array of their label values, or NULL if there are no labels.

The choose function can determine either that the new value matches one of the existing child nodes, or that a new child node must be added, or that the new value is inconsistent with the tuple prefix and so the inner tuple must be split to create a less restrictive prefix.

If the new value matches one of the existing child nodes, set resultType to spgMatchNode. Set nodeN to the index (from zero) of that node in the node array. Set levelAdd to the increment in level caused by descending through that node, or leave it as zero if the operator class does not use levels. Set restDatum to equal leafDatum if the operator class does not modify datums from one level to the next, or otherwise set it to the modified value to be used as leafDatum at the next level.

If a new child node must be added, set resultType to spgAddNode. Set nodeLabel to the label to be used for the new node, and set nodeN to the index (from zero) at which to insert the node in the node array. After the node has been added, the choose function will be called again with the modified inner tuple; that call should result in an spgMatchNode result.

If the new value is inconsistent with the tuple prefix, set resultType to spgSplitTuple. This action moves all the existing nodes into a new lower-level inner tuple, and replaces the existing inner tuple with a tuple having a single downlink pointing to the new lower-level inner tuple. Set prefixHasPrefix to indicate whether the new upper tuple should have a prefix, and if so set prefixPrefixDatum to the prefix value. This new prefix value must be sufficiently less restrictive than the original to accept the new value to be indexed. Set prefixNNodes to the number of nodes needed in the new tuple, and set prefixNodeLabels to a malloc'd array holding their labels, or to NULL if node labels are not required. Note that the total size of the new upper tuple must be no more than the total size of the tuple it is replacing; this constrains the lengths of the new prefix and new labels. Set childNodeN to the index (from zero) of the node that will downlink to the new lower-level inner tuple. Set postfixHasPrefix to indicate whether the new lower-level inner tuple should have a prefix, and if so set postfixPrefixDatum to the prefix value. The combination of these two prefixes and the downlink node's label (if any) must have the same meaning as the original prefix, because there is no opportunity to alter the node labels that are moved to the new lower-level tuple, nor to change any child index entries. After the node has been split, the choose function will be called again with the replacement inner tuple. That call may return an spgAddNode result, if no suitable node was created by the spgSplitTuple action. Eventually choose must return spgMatchNode to allow the insertion to descend to the next level.

picksplit

Decides how to create a new inner tuple over a set of leaf tuples.
The SQL declaration of the function must look like this:

```sql
CREATE FUNCTION my_picksplit(internal, internal) RETURNS void ...
```

The first argument is a pointer to a `spgPickSplitIn` C struct, containing input data for the function. The second argument is a pointer to a `spgPickSplitOut` C struct, which the function must fill with result data.

```c
typedef struct spgPickSplitIn {
    int nTuples;        /* number of leaf tuples */
    Datum *datums;         /* their datums (array of length nTuples) */
    int level;          /* current level (counting from zero) */
} spgPickSplitIn;

typedef struct spgPickSplitOut {
    bool hasPrefix;      /* new inner tuple should have a prefix */
    Datum prefixDatum;    /* if so, its value */
    int nNodes;         /* number of nodes for new inner tuple */
    Datum *nodeLabels;     /* their labels (or NULL for no labels) */
    int *mapTuplesToNodes; /* node index for each leaf tuple */
    Datum *leafTupleDatums; /* datum to store in each new leaf tuple */
} spgPickSplitOut;
```

The `nTuples` is the number of leaf tuples provided. `datums` is an array of their datum values of `spgConfigOut.leafType`. `level` is the current level that all the leaf tuples share, which will become the level of the new inner tuple.

Set `hasPrefix` to indicate whether the new inner tuple should have a prefix, and if so set `prefixDatum` to the prefix value. Set `nNodes` to indicate the number of nodes that the new inner tuple will contain, and set `nodeLabels` to an array of their label values, or to `NULL` if node labels are not required. Set `mapTuplesToNodes` to an array that gives the index (from zero) of the node that each leaf tuple should be assigned to. Set `leafTupleDatums` to an array of the values to be stored in the new leaf tuples (these will be the same as the input `datums` if the operator class does not modify datums from one level to the next). Note that the `picksplit` function is responsible for `palloc`ing the `nodeLabels`, `mapTuplesToNodes` and `leafTupleDatums` arrays.

If more than one leaf tuple is supplied, it is expected that the `picksplit` function will classify them into more than one node; otherwise it is not possible to split the leaf tuples across multiple pages, which is the ultimate purpose of this operation. Therefore, if the `picksplit` function ends up placing all the leaf tuples in the same node, the core SP-GiST code will override that decision and generate an inner tuple in which the leaf tuples are assigned at random to several identically-labeled nodes. Such a tuple is marked `allTheSame` to signify that this has happened. The choose and `inner_consistent` functions must take suitable care with such inner tuples. See Section 65.4.3 for more information.
The SQL declaration of the function must look like this:

```sql
CREATE FUNCTION my_inner_consistent(internal, internal) RETURNS void ...
```

The first argument is a pointer to a `spgInnerConsistentIn` C struct, containing input data for the function. The second argument is a pointer to a `spgInnerConsistentOut` C struct, which the function must fill with result data.

```c
typedef struct spgInnerConsistentIn {
    ScanKey     scankeys;       /* array of operators and comparison values */
    int         nkeys;          /* length of array */
    Datum       reconstructedValue;     /* value reconstructed at parent */
    void       *traversalValue; /* opclass-specific traverse value */
    MemoryContext traversalMemoryContext;   /* put new traverse values here */
    int         level;          /* current level (counting from zero) */
    bool        returnData;     /* original data must be returned? */
    bool        allTheSame;     /* tuple is marked all-the-same? */
    bool        hasPrefix;      /* tuple has a prefix? */
    Datum       prefixDatum;    /* if so, the prefix value */
    int         nNodes;         /* number of nodes in the inner tuple */
    Datum      *nodeLabels;     /* node label values (NULL if none) */
} spgInnerConsistentIn;

typedef struct spgInnerConsistentOut {
    int         nNodes;         /* number of child nodes to be visited */
    int         *nodeNumbers;   /* their indexes in the node array */
    int         *levelAdds;     /* increment level by this much for each */
    Datum      *reconstructedValues;    /* associated reconstructed values */
};
```

inner_consistent

Returns set of nodes (branches) to follow during tree search.
SP-GiST Indexes

void **traversalValues; /* opclass-specific
traverse values */

The array scankeys, of length nkeys, describes the index search condition(s). These conditions are combined with AND — only index entries that satisfy all of them are interesting. (Note that nkeys = 0 implies that all index entries satisfy the query.) Usually the consistent function only cares about the sk_strategy and sk_argument fields of each array entry, which respectively give the indexable operator and comparison value. In particular it is not necessary to check sk_flags to see if the comparison value is NULL, because the SP-GiST core code will filter out such conditions. reconstructedValue is the value reconstructed for the parent tuple; it is (Datum) 0 at the root level or if the inner_consistent function did not provide a value at the parent level. reconstructedValue is always of spgConfigOut.leafType type.

traversalValue is a pointer to any traverse data passed down from the previous call of inner_consistent on the parent index tuple, or NULL at the root level. traversalMemoryContext is the memory context in which to store output traverse values (see below). level is the current inner tuple's level, starting at zero for the root level. returnData is true if reconstructed data is required for this query; this will only be so if the config function asserted canReturnData. allTheSame is true if the current inner tuple is marked "all-the-same"; in this case all the nodes have the same label (if any) and so either all or none of them match the query (see Section 65.4.3). hasPrefix is true if the current inner tuple contains a prefix; if so, prefixDatum is its value. nNodes is the number of child nodes contained in the inner tuple, and nodeLabels is an array of their label values, or NULL if the nodes do not have labels.

nNodes must be set to the number of child nodes that need to be visited by the search, and nodeNumbers must be set to an array of their indexes. If the operator class keeps track of levels, set levelAdds to an array of the level increments required when descending to each node to be visited. (Often these increments will be the same for all the nodes, but that's not necessarily so, so an array is used.) If value reconstruction is needed, set reconstructedValues to an array of the values of spgConfigOut.leafType type reconstructed for each child node to be visited; otherwise, leave reconstructedValues as NULL. If it is desired to pass down additional out-of-band information ("traverse values") to lower levels of the tree search, set traversalValues to an array of the appropriate traverse values, one for each child node to be visited; otherwise, leave traversalValues as NULL. Note that the inner_consistent function is responsible for palloc'ing the nodeNumbers, levelAdds, reconstructedValues, and traversalValues arrays in the current memory context. However, any output traverse values pointed to by the traversalValues array should be allocated in traversalMemoryContext. Each traverse value must be a single palloc'd chunk.

leaf_consistent

Returns true if a leaf tuple satisfies a query.

The SQL declaration of the function must look like this:

CREATE FUNCTION my_leaf_consistent(internal, internal) RETURNS bool ...

The first argument is a pointer to a spgLeafConsistentIn C struct, containing input data for the function. The second argument is a pointer to a spgLeafConsistentOut C struct, which the function must fill with result data.

typedef struct spgLeafConsistentIn
{
    ScanKey     scankeys;   /* array of operators and
    int         nkeys;      /* length of array */

2226
The array scankeys, of length nkeys, describes the index search condition(s). These conditions are combined with AND — only index entries that satisfy all of them satisfy the query. (Note that nkeys = 0 implies that all index entries satisfy the query.) Usually the consistent function only cares about the sk_strategy and sk_argument fields of each array entry, which respectively give the indexable operator and comparison value. In particular it is not necessary to check sk_flags to see if the comparison value is NULL, because the SP-GiST core code will filter out such conditions. reconstructedValue is the value reconstructed for the parent tuple; it is (Datum) 0 at the root level or if the inner_consistent function did not provide a value at the parent level. reconstructedValue is always of spgConfigOut.leafType type. traversalValue is a pointer to any traverse data passed down from the previous call of inner_consistent on the parent index tuple, or NULL at the root level. level is the current leaf tuple’s level, starting at zero for the root level. returnData is true if reconstructed data is required for this query; this will only be so if the config function asserted canReturnData. leafDatum is the key value of spgConfigOut.leafType stored in the current leaf tuple.

The function must return true if the leaf tuple matches the query, or false if not. In the true case, if returnData is true then leafValue must be set to the value of spgConfigIn.attType originally supplied to be indexed for this leaf tuple. Also, recheck may be set to true if the match is uncertain and so the operator(s) must be re-applied to the actual heap tuple to verify the match.

The optional user-defined method is:

Datum compress(Datum in)

Converts the data item into a format suitable for physical storage in a leaf tuple of index page. It accepts spgConfigIn.attType value and return spgConfigOut.leafType value. Output value should not be toasted.

All the SP-GiST support methods are normally called in a short-lived memory context; that is, CurrentMemoryContext will be reset after processing of each tuple. It is therefore not very important to worry about pfree’ing everything you palloc. (The config method is an exception: it should try to avoid leaking memory. But usually the config method need do nothing but assign constants into the passed parameter struct.)

If the indexed column is of a collatable data type, the index collation will be passed to all the support methods, using the standard PG_GET_COLLATION() mechanism.
65.4. Implementation

This section covers implementation details and other tricks that are useful for implementers of SP-GiST operator classes to know.

65.4.1. SP-GiST Limits

Individual leaf tuples and inner tuples must fit on a single index page (8kB by default). Therefore, when indexing values of variable-length data types, long values can only be supported by methods such as radix trees, in which each level of the tree includes a prefix that is short enough to fit on a page, and the final leaf level includes a suffix also short enough to fit on a page. The operator class should set `longValuesOK` to true only if it is prepared to arrange for this to happen. Otherwise, the SP-GiST core will reject any request to index a value that is too large to fit on an index page.

Likewise, it is the operator class's responsibility that inner tuples do not grow too large to fit on an index page; this limits the number of child nodes that can be used in one inner tuple, as well as the maximum size of a prefix value.

Another limitation is that when an inner tuple's node points to a set of leaf tuples, those tuples must all be in the same index page. (This is a design decision to reduce seeking and save space in the links that chain such tuples together.) If the set of leaf tuples grows too large for a page, a split is performed and an intermediate inner tuple is inserted. For this to fix the problem, the new inner tuple must divide the set of leaf values into more than one node group. If the operator class's `picksplit` function fails to do that, the SP-GiST core resorts to extraordinary measures described in Section 65.4.3.

65.4.2. SP-GiST Without Node Labels

Some tree algorithms use a fixed set of nodes for each inner tuple; for example, in a quad-tree there are always exactly four nodes corresponding to the four quadrants around the inner tuple's centroid point. In such a case the code typically works with the nodes by number, and there is no need for explicit node labels. To suppress node labels (and thereby save some space), the `picksplit` function can return NULL for the `nodeLabels` array, and likewise the `choose` function can return NULL for the `prefixNodeLabels` array during a `spgSplitTuple` action. This will in turn result in `nodeLabels` being NULL during subsequent calls to `choose` and `inner_consistent`. In principle, node labels could be used for some inner tuples and omitted for others in the same index.

When working with an inner tuple having unlabeled nodes, it is an error for `choose` to return `spgAddNode`, since the set of nodes is supposed to be fixed in such cases.

65.4.3. “All-the-same” Inner Tuples

The SP-GiST core can override the results of the operator class's `picksplit` function when `picksplit` fails to divide the supplied leaf values into at least two node categories. When this happens, the new inner tuple is created with multiple nodes that each have the same label (if any) that `picksplit` gave to the one node it did use, and the leaf values are divided at random among these equivalent nodes. The `allTheSame` flag is set on the inner tuple to warn the `choose` and `inner_consistent` functions that the tuple does not have the node set that they might otherwise expect.

When dealing with an `allTheSame` tuple, a `choose` result of `spgMatchNode` is interpreted to mean that the new value can be assigned to any of the equivalent nodes; the core code will ignore the supplied `nodeN` value and descend into one of the nodes at random (so as to keep the tree balanced). It is an error for `choose` to return `spgAddNode`, since that would make the nodes not all equivalent; the `spgSplitTuple` action must be used if the value to be inserted doesn’t match the existing nodes.

When dealing with an `allTheSame` tuple, the `inner_consistent` function should return either all or none of the nodes as targets for continuing the index search, since they are all equivalent. This
may or may not require any special-case code, depending on how much the \texttt{inner\_consistent} function normally assumes about the meaning of the nodes.

### 65.5. Examples

The PostgreSQL source distribution includes several examples of index operator classes for SP-GiST, as described in Table 65.1. Look into \texttt{src/backend/access/spgist/} and \texttt{src/backend/utils/adt/} to see the code.
Chapter 66. GIN Indexes

66.1. Introduction

GIN stands for Generalized Inverted Index. GIN is designed for handling cases where the items to be indexed are composite values, and the queries to be handled by the index need to search for element values that appear within the composite items. For example, the items could be documents, and the queries could be searches for documents containing specific words.

We use the word *item* to refer to a composite value that is to be indexed, and the word *key* to refer to an element value. GIN always stores and searches for keys, not item values per se.

A GIN index stores a set of (key, posting list) pairs, where a *posting list* is a set of row IDs in which the key occurs. The same row ID can appear in multiple posting lists, since an item can contain more than one key. Each key value is stored only once, so a GIN index is very compact for cases where the same key appears many times.

GIN is generalized in the sense that the GIN access method code does not need to know the specific operations that it accelerates. Instead, it uses custom strategies defined for particular data types. The strategy defines how keys are extracted from indexed items and query conditions, and how to determine whether a row that contains some of the key values in a query actually satisfies the query.

One advantage of GIN is that it allows the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert. This is much the same advantage as using GiST.

The GIN implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov. There is more information about GIN on their website\(^1\).

66.2. Built-in Operator Classes

The core PostgreSQL distribution includes the GIN operator classes shown in Table 66.1. (Some of the optional modules described in Appendix F provide additional GIN operator classes.)

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>array_ops</td>
<td>anyarray</td>
<td>&amp;&amp; &lt;@ = @&gt;</td>
</tr>
<tr>
<td>jsonb_ops</td>
<td>jsonb</td>
<td>@? ?&amp; ?</td>
</tr>
<tr>
<td>jsonb_path_ops</td>
<td>jsonb</td>
<td>@&gt;</td>
</tr>
<tr>
<td>tsvector_ops</td>
<td>tsvector</td>
<td>@@@</td>
</tr>
</tbody>
</table>

Of the two operator classes for type *jsonb*, *jsonb_ops* is the default. *jsonb_path_ops* supports fewer operators but offers better performance for those operators. See Section 8.14.4 for details.

66.3. Extensibility

The GIN interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The GIN layer itself takes care of concurrency, logging and searching the tree structure.

\(^1\) http://www.sai.msu.su/~megera/wiki/Gin
All it takes to get a GIN access method working is to implement a few user-defined methods, which define the behavior of keys in the tree and the relationships between keys, indexed items, and indexable queries. In short, GIN combines extensibility with generality, code reuse, and a clean interface.

There are two methods that an operator class for GIN must provide:

Datum *extractValue(Datum itemValue, int32 *nkeys, bool **nullFlags)

Returns a palloc’d array of keys given an item to be indexed. The number of returned keys must be stored into *nkeys. If any of the keys can be null, also palloc an array of *nkeys bool fields, store its address at *nullFlags, and set these null flags as needed. *nullFlags can be left NULL (its initial value) if all keys are non-null. The return value can be NULL if the item contains no keys.

Datum *extractQuery(Datum query, int32 *nkeys, StrategyNumber n, bool **pmatch, Pointer **extra_data, bool **nullFlags, int32 *searchMode)

Returns a palloc’d array of keys given a value to be queried; that is, query is the value on the right-hand side of an indexable operator whose left-hand side is the indexed column. n is the strategy number of the operator within the operator class (see Section 38.15.2). Often, extractQuery will need to consult n to determine the data type of query and the method it should use to extract key values. The number of returned keys must be stored into *nkeys. If any of the keys can be null, also palloc an array of *nkeys bool fields, store its address at *nullFlags, and set these null flags as needed. *nullFlags can be left NULL (its initial value) if all keys are non-null. The return value can be NULL if the query contains no keys.

searchMode is an output argument that allows extractQuery to specify details about how the search will be done. If *searchMode is set to GIN_SEARCH_MODE_DEFAULT (which is the value it is initialized to before call), only items that match at least one of the returned keys are considered candidate matches. If *searchMode is set to GIN_SEARCH_MODE_INCLUDE_EMPTY, then in addition to items containing at least one matching key, items that contain no keys at all are considered candidate matches. (This mode is useful for implementing is-subset-of operators, for example.) If *searchMode is set to GIN_SEARCH_MODE_ALL, then all non-null items in the index are considered candidate matches, whether they match any of the returned keys or not. (This mode is much slower than the other two choices, since it requires scanning essentially the entire index, but it may be necessary to implement corner cases correctly. An operator that needs this mode in most cases is probably not a good candidate for a GIN operator class.) The symbols to use for setting this mode are defined in access/gin.h.

pmatch is an output argument for use when partial match is supported. To use it, extractQuery must allocate an array of *nkeys bools and store its address at *pmatch. Each element of the array should be set to true if the corresponding key requires partial match, false if not. If *pmatch is set to NULL then GIN assumes partial match is not required. The variable is initialized to NULL before call, so this argument can simply be ignored by operator classes that do not support partial match.

extra_data is an output argument that allows extractQuery to pass additional data to the consistent and comparePartial methods. To use it, extractQuery must allocate an array of *nkeys pointers and store its address at *extra_data, then store whatever it wants to into the individual pointers. The variable is initialized to NULL before call, so this argument can simply be ignored by operator classes that do not require extra data. If *extra_data is set, the whole array is passed to the consistent method, and the appropriate element to the comparePartial method.

An operator class must also provide a function to check if an indexed item matches the query. It comes in two flavors, a boolean consistent function, and a ternary triConsistent function. triConsistent covers the functionality of both, so providing triConsistent alone is sufficient. However, if the boolean variant is significantly cheaper to calculate, it can be advantageous to provide both. If only the boolean variant is provided, some optimizations that depend on refuting index items before fetching all the keys are disabled.
bool consistent(bool check[], StrategyNumber n, Datum query, int32 nkeys, Pointer extra_data[], bool *recheck, Datum queryKeys[], bool nullFlags[])  

Returns true if an indexed item satisfies the query operator with strategy number n (or might satisfy it, if the recheck indication is returned). This function does not have direct access to the indexed item's value, since GIN does not store items explicitly. Rather, what is available is knowledge about which key values extracted from the query appear in a given indexed item. The check array has length nkeys, which is the same as the number of keys previously returned by extractQuery for this query datum. Each element of the check array is true if the indexed item contains the corresponding query key, i.e., if (check[i] == true) the i-th key of the extractQuery result array is present in the indexed item. The original query datum is passed in case the consistent method needs to consult it, and so are the queryKeys[] and nullFlags[] arrays previously returned by extractQuery. extra_data is the extra-data array returned by extractQuery, or NULL if none.

When extractQuery returns a null key in queryKeys[], the corresponding check[] element is true if the indexed item contains a null key; that is, the semantics of check[] are like IS NOT DISTINCT FROM. The consistent function can examine the corresponding null-Flags[] element if it needs to tell the difference between a regular value match and a null match.

On success, *recheck should be set to true if the heap tuple needs to be rechecked against the query operator, or false if the index test is exact. That is, a false return value guarantees that the heap tuple does not match the query; a true return value with *recheck set to false guarantees that the heap tuple does match the query; and a true return value with *recheck set to true means that the heap tuple might match the query, so it needs to be fetched and rechecked by evaluating the query operator directly against the originally indexed item.

GinTernaryValue triConsistent(GinTernaryValue check[], StrategyNumber n, Datum query, int32 nkeys, Pointer extra_data[], Datum queryKeys[], bool nullFlags[])  

triConsistent is similar to consistent, but instead of booleans in the check vector, there are three possible values for each key: GIN_TRUE, GIN_FALSE and GIN_MAYBE. GIN_FALSE and GIN_TRUE have the same meaning as regular boolean values, while GIN_MAYBE means that the presence of that key is not known. When GIN_MAYBE values are present, the function should only return GIN_TRUE if the item certainly matches whether or not the index item contains the corresponding query keys. Likewise, the function must return GIN_FALSE only if the item certainly does not match, whether or not it contains the GIN_MAYBE keys. If the result depends on the GIN_MAYBE entries, i.e., the match cannot be confirmed or refuted based on the known query keys, the function must return GIN_MAYBE.

When there are no GIN_MAYBE values in the check vector, a GIN_MAYBE return value is the equivalent of setting the recheck flag in the boolean consistent function.

In addition, GIN must have a way to sort the key values stored in the index. The operator class can define the sort ordering by specifying a comparison method:

int compare(Datum a, Datum b)

Compares two keys (not indexed items!) and returns an integer less than zero, zero, or greater than zero, indicating whether the first key is less than, equal to, or greater than the second. Null keys are never passed to this function.

Alternatively, if the operator class does not provide a compare method, GIN will look up the default btree operator class for the index key data type, and use its comparison function. It is recommended to specify the comparison function in a GIN operator class that is meant for just one data type, as looking up the btree operator class costs a few cycles. However, polymorphic GIN operator classes (such as array_ops) typically cannot specify a single comparison function.
Optionally, an operator class for GIN can supply the following method:

```c
int comparePartial(Datum partial_key, Datum key, StrategyNumber n, Pointer extra_data)
```

Compare a partial-match query key to an index key. Returns an integer whose sign indicates the result: less than zero means the index key does not match the query, but the index scan should continue; zero means that the index key does match the query; greater than zero indicates that the index scan should stop because no more matches are possible. The strategy number \( n \) of the operator that generated the partial match query is provided, in case its semantics are needed to determine when to end the scan. Also, \( \text{extra_data} \) is the corresponding element of the extra-data array made by \text{extractQuery}, or \text{NULL} if none. Null keys are never passed to this function.

To support “partial match” queries, an operator class must provide the \text{comparePartial} method, and its \text{extractQuery} method must set the \text{pmatch} parameter when a partial-match query is encountered. See Section 66.4.2 for details.

The actual data types of the various \text{Datum} values mentioned above vary depending on the operator class. The item values passed to \text{extractValue} are always of the operator class's input type, and all key values must be of the class's \text{STORAGE} type. The type of the \text{query} argument passed to \text{extractQuery}, \text{consistent} and \text{triConsistent} is whatever is the right-hand input type of the class member operator identified by the strategy number. This need not be the same as the indexed type, so long as key values of the correct type can be extracted from it. However, it is recommended that the SQL declarations of these three support functions use the opclass's indexed data type for the \text{query} argument, even though the actual type might be something else depending on the operator.

### 66.4. Implementation

Internally, a GIN index contains a B-tree index constructed over keys, where each key is an element of one or more indexed items (a member of an array, for example) and where each tuple in a leaf page contains either a pointer to a B-tree of heap pointers (a “posting tree”), or a simple list of heap pointers (a “posting list”) when the list is small enough to fit into a single index tuple along with the key value.

As of PostgreSQL 9.1, null key values can be included in the index. Also, placeholder nulls are included in the index for indexed items that are null or contain no keys according to \text{extractValue}. This allows searches that should find empty items to do so.

Multicolumn GIN indexes are implemented by building a single B-tree over composite values (column number, key value). The key values for different columns can be of different types.

#### 66.4.1. GIN Fast Update Technique

Updating a GIN index tends to be slow because of the intrinsic nature of inverted indexes: inserting or updating one heap row can cause many inserts into the index (one for each key extracted from the indexed item). As of PostgreSQL 8.4, GIN is capable of postponing much of this work by inserting new tuples into a temporary, unsorted list of pending entries. When the table is vacuumed or autoanalyzed, or when \text{gin_clean_pending_list} function is called, or if the pending list becomes larger than \text{gin_pending_list_limit}, the entries are moved to the main GIN data structure using the same bulk insert techniques used during initial index creation. This greatly improves GIN index update speed, even counting the additional vacuum overhead. Moreover the overhead work can be done by a background process instead of in foreground query processing.

The main disadvantage of this approach is that searches must scan the list of pending entries in addition to searching the regular index, and so a large list of pending entries will slow searches significantly. Another disadvantage is that, while most updates are fast, an update that causes the pending list to become “too large” will incur an immediate cleanup cycle and thus be much slower than other updates. Proper use of autovacuum can minimize both of these problems.
66.4.2. Partial Match Algorithm

GIN can support “partial match” queries, in which the query does not determine an exact match for one or more keys, but the possible matches fall within a reasonably narrow range of key values (within the key sorting order determined by the compare support method). The extractQuery method, instead of returning a key value to be matched exactly, returns a key value that is the lower bound of the range to be searched, and sets the pmatch flag true. The key range is then scanned using the comparePartial method. comparePartial must return zero for a matching index key, less than zero for a non-match that is still within the range to be searched, or greater than zero if the index key is past the range that could match.

66.5. GIN Tips and Tricks

Create vs. insert

Insertion into a GIN index can be slow due to the likelihood of many keys being inserted for each item. So, for bulk insertions into a table it is advisable to drop the GIN index and recreate it after finishing bulk insertion.

As of PostgreSQL 8.4, this advice is less necessary since delayed indexing is used (see Section 66.4.1 for details). But for very large updates it may still be best to drop and recreate the index.

maintenance_work_mem

Build time for a GIN index is very sensitive to the maintenance_work_mem setting; it doesn’t pay to skimp on work memory during index creation.

gin_pending_list_limit

During a series of insertions into an existing GIN index that has fastupdate enabled, the system will clean up the pending-entry list whenever the list grows larger than gin_pending_list_limit. To avoid fluctuations in observed response time, it’s desirable to have pending-list cleanup occur in the background (i.e., via autovacuum). Foreground cleanup operations can be avoided by increasing gin_pending_list_limit or making autovacuum more aggressive. However, enlarging the threshold of the cleanup operation means that if a foreground cleanup does occur, it will take even longer.

gin_pending_list_limit can be overridden for individual GIN indexes by changing storage parameters, and which allows each GIN index to have its own cleanup threshold. For example, it’s possible to increase the threshold only for the GIN index which can be updated heavily, and decrease it otherwise.

gin_fuzzy_search_limit

The primary goal of developing GIN indexes was to create support for highly scalable full-text search in PostgreSQL, and there are often situations when a full-text search returns a very large set of results. Moreover, this often happens when the query contains very frequent words, so that the large result set is not even useful. Since reading many tuples from the disk and sorting them could take a lot of time, this is unacceptable for production. (Note that the index search itself is very fast.)

To facilitate controlled execution of such queries, GIN has a configurable soft upper limit on the number of rows returned: the gin_fuzzy_search_limit configuration parameter. It is set to 0 (meaning no limit) by default. If a non-zero limit is set, then the returned set is a subset of the whole result set, chosen at random.
"Soft" means that the actual number of returned results could differ somewhat from the specified limit, depending on the query and the quality of the system's random number generator.

From experience, values in the thousands (e.g., 5000 — 20000) work well.

### 66.6. Limitations

GIN assumes that indexable operators are strict. This means that `extractValue` will not be called at all on a null item value (instead, a placeholder index entry is created automatically), and `extractQuery` will not be called on a null query value either (instead, the query is presumed to be unsatisfiable). Note however that null key values contained within a non-null composite item or query value are supported.

### 66.7. Examples

The core PostgreSQL distribution includes the GIN operator classes previously shown in Table 66.1. The following `contrib` modules also contain GIN operator classes:

- `btree_gin`
  
  B-tree equivalent functionality for several data types

- `hstore`
  
  Module for storing (key, value) pairs

- `intarray`
  
  Enhanced support for `int[]`

- `pg_trgm`
  
  Text similarity using trigram matching
Chapter 67. BRIN Indexes

67.1. Introduction

BRIN stands for Block Range Index. BRIN is designed for handling very large tables in which certain columns have some natural correlation with their physical location within the table. A block range is a group of pages that are physically adjacent in the table; for each block range, some summary info is stored by the index. For example, a table storing a store’s sale orders might have a date column on which each order was placed, and most of the time the entries for earlier orders will appear earlier in the table as well; a table storing a ZIP code column might have all codes for a city grouped together naturally.

BRIN indexes can satisfy queries via regular bitmap index scans, and will return all tuples in all pages within each range if the summary info stored by the index is consistent with the query conditions. The query executor is in charge of rechecking these tuples and discarding those that do not match the query conditions — in other words, these indexes are lossy. Because a BRIN index is very small, scanning the index adds little overhead compared to a sequential scan, but may avoid scanning large parts of the table that are known not to contain matching tuples.

The specific data that a BRIN index will store, as well as the specific queries that the index will be able to satisfy, depend on the operator class selected for each column of the index. Data types having a linear sort order can have operator classes that store the minimum and maximum value within each block range, for instance; geometrical types might store the bounding box for all the objects in the block range.

The size of the block range is determined at index creation time by the pages_per_range storage parameter. The number of index entries will be equal to the size of the relation in pages divided by the selected value for pages_per_range. Therefore, the smaller the number, the larger the index becomes (because of the need to store more index entries), but at the same time the summary data stored can be more precise and more data blocks can be skipped during an index scan.

67.1.1. Index Maintenance

At the time of creation, all existing heap pages are scanned and a summary index tuple is created for each range, including the possibly-incomplete range at the end. As new pages are filled with data, page ranges that are already summarized will cause the summary information to be updated with data from the new tuples. When a new page is created that does not fall within the last summarized range, that range does not automatically acquire a summary tuple; those tuples remain unsummarized until a summarization run is invoked later, creating initial summaries. This process can be invoked manually using the brin_summarize_range(regclass, bigint) or brin_summarize_new_values(regclass) functions; automatically when VACUUM processes the table; or by automatic summarization executed by autovacuum, as insertions occur. (This last trigger is disabled by default and can be enabled with the autosummarize parameter.) Conversely, a range can be de-summarized using the brin_desummarize_range(regclass, bigint) function, which is useful when the index tuple is no longer a very good representation because the existing values have changed.

When autosummarization is enabled, each time a page range is filled a request is sent to autovacuum for it to execute a targeted summarization for that range, to be fulfilled at the end of the next worker run on the same database. If the request queue is full, the request is not recorded and a message is sent to the server log:

LOG: request for BRIN range summarization for index "brin_wi_idx" page 128 was not recorded

When this happens, the range will be summarized normally during the next regular vacuum of the table.
67.2. Built-in Operator Classes

The core PostgreSQL distribution includes the BRIN operator classes shown in Table 67.1.

The minmax operator classes store the minimum and the maximum values appearing in the indexed column within the range. The inclusion operator classes store a value which includes the values in the indexed column within the range.

### Table 67.1. Built-in BRIN Operator Classes

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>abstime_minmax_ops</td>
<td>abstime</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>int8_minmax_ops</td>
<td>bigint</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>bit_minmax_ops</td>
<td>bit</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>varbit_minmax_ops</td>
<td>bit varying</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>box_inclusion_ops</td>
<td>box</td>
<td>&lt;&lt; &amp; &amp; &amp; &gt;&gt; ~ = &gt; @ &amp; &lt;</td>
</tr>
<tr>
<td>bytea_minmax_ops</td>
<td>bytea</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>bpchar_minmax_ops</td>
<td>character</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>char_minmax_ops</td>
<td>&quot;char&quot;</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>date_minmax_ops</td>
<td>date</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>float8_minmax_ops</td>
<td>double precision</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>inet_minmax_ops</td>
<td>inet</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>network_inclusion_ops</td>
<td>inet</td>
<td>&amp; &amp; &gt;= &lt; &lt;= &gt; &gt; &lt;&lt;</td>
</tr>
<tr>
<td>int4_minmax_ops</td>
<td>integer</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>interval_minmax_ops</td>
<td>interval</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>macaddr_minmax_ops</td>
<td>macaddr</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>macaddr8_minmax_ops</td>
<td>macaddr8</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>name_minmax_ops</td>
<td>name</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>numeric_minmax_ops</td>
<td>numeric</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>pg_lsn_minmax_ops</td>
<td>pg_lsn</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>oid_minmax_ops</td>
<td>oid</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>range_inclusion_ops</td>
<td>any range type</td>
<td>&lt;&lt; &amp; &amp; &amp; &amp; &gt;&gt; @ &gt; @</td>
</tr>
<tr>
<td>float4_minmax_ops</td>
<td>real</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>reltime_minmax_ops</td>
<td>reltime</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>int2_minmax_ops</td>
<td>smallint</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>text_minmax_ops</td>
<td>text</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>tid_minmax_ops</td>
<td>tid</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>timestamp_minmax_ops</td>
<td>timestamp without time zone</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>timestamptz_minmax_ops</td>
<td>timestamp with time zone</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>time_minmax_ops</td>
<td>time without time zone</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>timetz_minmax_ops</td>
<td>time with time zone</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
<tr>
<td>uuid_minmax_ops</td>
<td>uuid</td>
<td>&lt; &lt;= = &gt;= &gt;</td>
</tr>
</tbody>
</table>
67.3. Extensibility

The BRIN interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The BRIN layer itself takes care of concurrency, logging and searching the index structure.

All it takes to get a BRIN access method working is to implement a few user-defined methods, which define the behavior of summary values stored in the index and the way they interact with scan keys. In short, BRIN combines extensibility with generality, code reuse, and a clean interface.

There are four methods that an operator class for BRIN must provide:

BrinOpcInfo *opcInfo(Oid type_oid)

Returns internal information about the indexed columns' summary data. The return value must point to a palloc'd BrinOpcInfo, which has this definition:

typedef struct BrinOpcInfo
{
    /* Number of columns stored in an index column of this opclass */
    uint16      oi_nstored;

    /* Opaque pointer for the opclass' private use */
    void       *oi_opaque;

    /* Type cache entries of the stored columns */
    TypeCacheEntry *oi_typcache[ FLEXIBLE_ARRAY_MEMBER ];
} BrinOpcInfo;

BrinOpcInfo.oopaque can be used by the operator class routines to pass information between support functions during an index scan.

bool consistent(BrinDesc *bdesc, BrinValues *column, ScanKey key)

Returns whether the ScanKey is consistent with the given indexed values for a range. The attribute number to use is passed as part of the scan key.

bool addValue(BrinDesc *bdesc, BrinValues *column, Datum newval, bool isnull)

Given an index tuple and an indexed value, modifies the indicated attribute of the tuple so that it additionally represents the new value. If any modification was done to the tuple, true is returned.

bool unionTuples(BrinDesc *bdesc, BrinValues *a, BrinValues *b)

Consolidates two index tuples. Given two index tuples, modifies the indicated attribute of the first of them so that it represents both tuples. The second tuple is not modified.

The core distribution includes support for two types of operator classes: minmax and inclusion. Operator class definitions using them are shipped for in-core data types as appropriate. Additional operator classes can be defined by the user for other data types using equivalent definitions, without having to write any source code; appropriate catalog entries being declared is enough. Note that assumptions about the semantics of operator strategies are embedded in the support functions' source code.

Operator classes that implement completely different semantics are also possible, provided implementations of the four main support functions described above are written. Note that backwards compat-
ibility across major releases is not guaranteed: for example, additional support functions might be required in later releases.

To write an operator class for a data type that implements a totally ordered set, it is possible to use the minmax support functions alongside the corresponding operators, as shown in Table 67.2. All operator class members (functions and operators) are mandatory.

**Table 67.2. Function and Support Numbers for Minmax Operator Classes**

<table>
<thead>
<tr>
<th>Operator class member</th>
<th>Object</th>
</tr>
</thead>
<tbody>
<tr>
<td>Support Function 1</td>
<td>internal function brin_minmax_opcinfo()</td>
</tr>
<tr>
<td>Support Function 2</td>
<td>internal function brin_minmax_add_value()</td>
</tr>
<tr>
<td>Support Function 3</td>
<td>internal function brin_minmax_consistent()</td>
</tr>
<tr>
<td>Support Function 4</td>
<td>internal function brin_minmax_union()</td>
</tr>
<tr>
<td>Operator Strategy 1</td>
<td>operator less-than</td>
</tr>
<tr>
<td>Operator Strategy 2</td>
<td>operator less-than-or-equal-to</td>
</tr>
<tr>
<td>Operator Strategy 3</td>
<td>operator equal-to</td>
</tr>
<tr>
<td>Operator Strategy 4</td>
<td>operator greater-than-or-equal-to</td>
</tr>
<tr>
<td>Operator Strategy 5</td>
<td>operator greater-than</td>
</tr>
</tbody>
</table>

To write an operator class for a complex data type which has values included within another type, it's possible to use the inclusion support functions alongside the corresponding operators, as shown in Table 67.3. It requires only a single additional function, which can be written in any language. More functions can be defined for additional functionality. All operators are optional. Some operators require other operators, as shown as dependencies on the table.

**Table 67.3. Function and Support Numbers for Inclusion Operator Classes**

<table>
<thead>
<tr>
<th>Operator class member</th>
<th>Object</th>
<th>Dependency</th>
</tr>
</thead>
<tbody>
<tr>
<td>Support Function 1</td>
<td>internal function brin_inclusion_opcinfo()</td>
<td></td>
</tr>
<tr>
<td>Support Function 2</td>
<td>internal function brin_inclusion_add_value()</td>
<td></td>
</tr>
<tr>
<td>Support Function 3</td>
<td>internal function brin_inclusion_consistent()</td>
<td></td>
</tr>
<tr>
<td>Support Function 4</td>
<td>internal function brin_inclusion_union()</td>
<td></td>
</tr>
<tr>
<td>Support Function 11</td>
<td>function to merge two elements</td>
<td></td>
</tr>
<tr>
<td>Support Function 12</td>
<td>optional function to check whether two elements are mergeable</td>
<td>Operator Strategy 4</td>
</tr>
<tr>
<td>Support Function 13</td>
<td>optional function to check if an element is contained within another</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>Support Function 14</td>
<td>optional function to check whether an element is empty</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>Operator Strategy 1</td>
<td>operator left-of</td>
<td>Operator Strategy 4</td>
</tr>
</tbody>
</table>

2239
### Operator class member

<table>
<thead>
<tr>
<th>Operator Strategy</th>
<th>Object</th>
<th>Dependency</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>operator overlaps</td>
<td>Operator Strategy 1</td>
</tr>
<tr>
<td>4</td>
<td>operator does-not-extend-to-the-left-of</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>operator right-of</td>
<td>Operator Strategy 2</td>
</tr>
<tr>
<td>6, 18</td>
<td>operator same-as-or-equal-to</td>
<td>Operator Strategy 7</td>
</tr>
<tr>
<td>7, 13, 16, 24, 25</td>
<td>operator contains-or-equal-to</td>
<td></td>
</tr>
<tr>
<td>8, 14, 26, 27</td>
<td>operator is-contained-by-or-equal-to</td>
<td>Operator Strategy 3</td>
</tr>
<tr>
<td>9</td>
<td>operator does-not-extend-above</td>
<td>Operator Strategy 11</td>
</tr>
<tr>
<td>10</td>
<td>operator is-below</td>
<td>Operator Strategy 12</td>
</tr>
<tr>
<td>11</td>
<td>operator is-above</td>
<td>Operator Strategy 9</td>
</tr>
<tr>
<td>12</td>
<td>operator does-not-extend-below</td>
<td>Operator Strategy 10</td>
</tr>
<tr>
<td>20</td>
<td>operator less-than</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>21</td>
<td>operator less-than-or-equal-to</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>22</td>
<td>operator greater-than</td>
<td>Operator Strategy 1</td>
</tr>
<tr>
<td>23</td>
<td>operator greater-than-or-equal-to</td>
<td>Operator Strategy 1</td>
</tr>
</tbody>
</table>

Support function numbers 1-10 are reserved for the BRIN internal functions, so the SQL level functions start with number 11. Support function number 11 is the main function required to build the index. It should accept two arguments with the same data type as the operator class, and return the union of them. The inclusion operator class can store union values with different data types if it is defined with the `STORAGE` parameter. The return value of the union function should match the `STORAGE` data type.

Support function numbers 12 and 14 are provided to support irregularities of built-in data types. Function number 12 is used to support network addresses from different families which are not mergeable. Function number 14 is used to support empty ranges. Function number 13 is an optional but recommended one, which allows the new value to be checked before it is passed to the union function. As the BRIN framework can shortcut some operations when the union is not changed, using this function can improve index performance.

Both minmax and inclusion operator classes support cross-data-type operators, though with these the dependencies become more complicated. The minmax operator class requires a full set of operators to be defined with both arguments having the same data type. It allows additional data types to be supported by defining extra sets of operators. Inclusion operator class operator strategies are dependent on another operator strategy as shown in Table 67.3, or the same operator strategy as themselves. They require the dependency operator to be defined with the `STORAGE` data type as the left-hand-side argument and the other supported data type to be the right-hand-side argument of the supported operator. See `float4_minmax_ops` as an example of minmax, and `box_inclusion_ops` as an example of inclusion.
Chapter 68. Database Physical Storage

This chapter provides an overview of the physical storage format used by PostgreSQL databases.

68.1. Database File Layout

This section describes the storage format at the level of files and directories.

Traditionally, the configuration and data files used by a database cluster are stored together within the cluster's data directory, commonly referred to as PGDATA (after the name of the environment variable that can be used to define it). A common location for PGDATA is /var/lib/pgsql/data. Multiple clusters, managed by different server instances, can exist on the same machine.

The PGDATA directory contains several subdirectories and control files, as shown in Table 68.1. In addition to these required items, the cluster configuration files postgresql.conf, pg_hba.conf, and pg_ident.conf are traditionally stored in PGDATA, although it is possible to place them elsewhere.

Table 68.1. Contents of PGDATA

<table>
<thead>
<tr>
<th>Item</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PG_VERSION</td>
<td>A file containing the major version number of PostgreSQL</td>
</tr>
<tr>
<td>base</td>
<td>Subdirectory containing per-database subdirectories</td>
</tr>
<tr>
<td>current_logfiles</td>
<td>File recording the log file(s) currently written to by the logging collector</td>
</tr>
<tr>
<td>global</td>
<td>Subdirectory containing cluster-wide tables, such as pg_database</td>
</tr>
<tr>
<td>pg_commit_ts</td>
<td>Subdirectory containing transaction commit timestamp data</td>
</tr>
<tr>
<td>pg_dynshmem</td>
<td>Subdirectory containing files used by the dynamic shared memory subsystem</td>
</tr>
<tr>
<td>pg_logical</td>
<td>Subdirectory containing status data for logical decoding</td>
</tr>
<tr>
<td>pg_multixact</td>
<td>Subdirectory containing multitransaction status data (used for shared row locks)</td>
</tr>
<tr>
<td>pg_notify</td>
<td>Subdirectory containing LISTEN/NOTIFY status data</td>
</tr>
<tr>
<td>pg_repislot</td>
<td>Subdirectory containing replication slot data</td>
</tr>
<tr>
<td>pg_serial</td>
<td>Subdirectory containing information about committed serializable transactions</td>
</tr>
<tr>
<td>pg_snapshots</td>
<td>Subdirectory containing exported snapshots</td>
</tr>
<tr>
<td>pg_stat</td>
<td>Subdirectory containing permanent files for the statistics subsystem</td>
</tr>
<tr>
<td>pg_stat_tmp</td>
<td>Subdirectory containing temporary files for the statistics subsystem</td>
</tr>
<tr>
<td>pg_subtrans</td>
<td>Subdirectory containing subtransaction status data</td>
</tr>
<tr>
<td>Item</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>pg_tblspc</td>
<td>Subdirectory containing symbolic links to tablespaces</td>
</tr>
<tr>
<td>pg_twophase</td>
<td>Subdirectory containing state files for prepared transactions</td>
</tr>
<tr>
<td>pg_wal</td>
<td>Subdirectory containing WAL (Write Ahead Log) files</td>
</tr>
<tr>
<td>pg_xact</td>
<td>Subdirectory containing transaction commit status data</td>
</tr>
<tr>
<td>postgresql.auto.conf</td>
<td>A file used for storing configuration parameters that are set by ALTER SYSTEM</td>
</tr>
<tr>
<td>postmaster.opts</td>
<td>A file recording the command-line options the server was last started with</td>
</tr>
<tr>
<td>postmaster.pid</td>
<td>A lock file recording the current postmaster process ID (PID), cluster data directory path, postmaster start timestamp, port number, Unix-domain socket directory path (empty on Windows), first valid listen_address (IP address or *, or empty if not listening on TCP), and shared memory segment ID (this file is not present after server shutdown)</td>
</tr>
</tbody>
</table>

For each database in the cluster there is a subdirectory within PGDATA/base, named after the database's OID in pg_database. This subdirectory is the default location for the database's files; in particular, its system catalogs are stored there.

Each table and index is stored in a separate file. For ordinary relations, these files are named after the table or index's filenode number, which can be found in pg_class.relfilenode. But for temporary relations, the file name is of the form tBBB_FFF, where BBB is the backend ID of the backend which created the file, and FFF is the filenode number. In either case, in addition to the main file (a/k/a main fork), each table and index has a free space map (see Section 68.3), which stores information about free space available in the relation. The free space map is stored in a file named with the filenode number plus the suffix _fsm. Tables also have a visibility map, stored in a fork with the suffix _vm, to track which pages are known to have no dead tuples. The visibility map is described further in Section 68.4. Unlogged tables and indexes have a third fork, known as the initialization fork, which is stored in a fork with the suffix _init (see Section 68.5).

**Caution**

Note that while a table's filenode often matches its OID, this is not necessarily the case; some operations, like TRUNCATE, REINDEX, CLUSTER and some forms of ALTER TABLE, can change the filenode while preserving the OID. Avoid assuming that filenode and table OID are the same. Also, for certain system catalogs including pg_class itself, pg_class.relfilenode contains zero. The actual filenode number of these catalogs is stored in a lower-level data structure, and can be obtained using the pg_relation_filenode() function.

When a table or index exceeds 1 GB, it is divided into gigabyte-sized segments. The first segment's file name is the same as the filenode; subsequent segments are named filenode.1, filenode.2, etc. This arrangement avoids problems on platforms that have file size limitations. (Actually, 1 GB is just the default segment size. The segment size can be adjusted using the configuration option --with-segsize when building PostgreSQL.) In principle, free space map and visibility map forks could require multiple segments as well, though this is unlikely to happen in practice.
A table that has columns with potentially large entries will have an associated TOAST table, which is used for out-of-line storage of field values that are too large to keep in the table rows proper. `pg_class.reltoastrelid` links from a table to its TOAST table, if any. See Section 68.2 for more information.

The contents of tables and indexes are discussed further in Section 68.6.

Tablespaces make the scenario more complicated. Each user-defined tablespace has a symbolic link inside the `PGDATA/pg_tblspc` directory, which points to the physical tablespace directory (i.e., the location specified in the tablespace’s `CREATE TABLESPACE` command). This symbolic link is named after the tablespace’s OID. Inside the physical tablespace directory there is a subdirectory with a name that depends on the PostgreSQL server version, such as `PG_9.0_20100805`. (The reason for using this subdirectory is so that successive versions of the database can use the same `CREATE TABLESPACE` location value without conflicts.) Within the version-specific subdirectory, there is a subdirectory for each database that has elements in the tablespace, named after the database’s OID. Tables and indexes are stored within that directory, using the filenode naming scheme. The `pg_default` tablespace is not accessed through `pg_tblspc`, but corresponds to `PGDATA/base`. Similarly, the `pg_global` tablespace is not accessed through `pg_tblspc`, but corresponds to `PGDATA/global`.

The `pg_relation_filepath()` function shows the entire path (relative to `PGDATA`) of any relation. It is often useful as a substitute for remembering many of the above rules. But keep in mind that this function just gives the name of the first segment of the main fork of the relation — you may need to append a segment number and/or `_fsm`, `_vm`, or `_init` to find all the files associated with the relation.

Temporary files (for operations such as sorting more data than can fit in memory) are created within `PGDATA/base/pgsql_tmp`, or within a `pgsql_tmp` subdirectory of a tablespace directory if a tablespace other than `pg_default` is specified for them. The name of a temporary file has the form `pgsql_tmpPPP.NNN`, where `PPP` is the PID of the owning backend and `NNN` distinguishes different temporary files of that backend.

## 68.2. TOAST

This section provides an overview of TOAST (The Oversized-Attribute Storage Technique).

PostgreSQL uses a fixed page size (commonly 8 kB), and does not allow tuples to span multiple pages. Therefore, it is not possible to store very large field values directly. To overcome this limitation, large field values are compressed and/or broken up into multiple physical rows. This happens transparently to the user, with only small impact on most of the backend code. The technique is affectionately known as TOAST (or “the best thing since sliced bread”). The TOAST infrastructure is also used to improve handling of large data values in-memory.

Only certain data types support TOAST — there is no need to impose the overhead on data types that cannot produce large field values. To support TOAST, a data type must have a variable-length (`varlena`) representation, in which, ordinarily, the first four-byte word of any stored value contains the total length of the value in bytes (including itself). TOAST does not constrain the rest of the data type’s representation. The special representations collectively called `TOASTed values` work by modifying or reinterpreting this initial length word. Therefore, the C-level functions supporting a TOAST-able data type must be careful about how they handle potentially `TOASTed` input values: an input might not actually consist of a four-byte length word and contents until after it’s been `detoasted`. (This is normally done by invoking `PG_DETOAST_DATUM` before doing anything with an input value, but in some cases more efficient approaches are possible. See Section 38.12.1 for more detail.)

TOAST usurps two bits of the varlena length word (the high-order bits on big-endian machines, the low-order bits on little-endian machines), thereby limiting the logical size of any value of a TOAST-able data type to 1 GB ($2^{30}$ - 1 bytes). When both bits are zero, the value is an ordinary un-`TOASTed` value of the data type, and the remaining bits of the length word give the total datum size (including length word) in bytes. When the highest-order or lowest-order bit is set, the value has only a single-byte
header instead of the normal four-byte header, and the remaining bits of that byte give the total datum size (including length byte) in bytes. This alternative supports space-efficient storage of values shorter than 127 bytes, while still allowing the data type to grow to 1 GB at need. Values with single-byte headers aren’t aligned on any particular boundary, whereas values with four-byte headers are aligned on at least a four-byte boundary; this omission of alignment padding provides additional space savings that is significant compared to short values. As a special case, if the remaining bits of a single-byte header are all zero (which would be impossible for a self-inclusive length), the value is a pointer to out-of-line data, with several possible alternatives as described below. The type and size of such a TOAST pointer are determined by a code stored in the second byte of the datum. Lastly, when the highest-order or lowest-order bit is clear but the adjacent bit is set, the content of the datum has been compressed and must be decompressed before use. In this case the remaining bits of the four-byte length word give the total size of the compressed datum, not the original data. Note that compression is also possible for out-of-line data but the varlena header does not tell whether it has occurred — the content of the TOAST pointer tells that, instead.

As mentioned, there are multiple types of TOAST pointer datums. The oldest and most common type is a pointer to out-of-line data stored in a TOAST table that is separate from, but associated with, the table containing the TOAST pointer datum itself. These on-disk pointer datums are created by the TOAST management code (in access/heap/tuptoaster.c) when a tuple to be stored on disk is too large to be stored as-is. Further details appear in Section 68.2.1. Alternatively, a TOAST pointer datum can contain a pointer to out-of-line data that appears elsewhere in memory. Such datums are necessarily short-lived, and will never appear on-disk, but they are very useful for avoiding copying and redundant processing of large data values. Further details appear in Section 68.2.2.

The compression technique used for either in-line or out-of-line compressed data is a fairly simple and very fast member of the LZ family of compression techniques. See src/common/pg_lzcompress.c for the details.

### 68.2.1. Out-of-line, on-disk TOAST storage

If any of the columns of a table are TOAST-able, the table will have an associated TOAST table, whose OID is stored in the table's pg_class.reltoastrelid entry. On-disk TOASTed values are kept in the TOAST table, as described in more detail below.

Out-of-line values are divided (after compression if used) into chunks of at most TOAST_MAX_CHUNK_SIZE bytes (by default this value is chosen so that four chunk rows will fit on a page, making it about 2000 bytes). Each chunk is stored as a separate row in the TOAST table belonging to the owning table. Every TOAST table has the columns chunk_id (an OID identifying the particular TOASTed value), chunk_seq (a sequence number for the chunk within its value), and chunk_data (the actual data of the chunk). A unique index on chunk_id and chunk_seq provides fast retrieval of the values. A pointer datum representing an out-of-line on-disk TOASTed value therefore needs to store the OID of the TOAST table in which to look and the OID of the specific value (its chunk_id). For convenience, pointer datums also store the logical datum size (original uncompressed data length) and physical stored size (different if compression was applied). Allowing for the varlena header bytes, the total size of an on-disk TOAST pointer datum is therefore 18 bytes regardless of the actual size of the represented value.

The TOAST management code is triggered only when a row value to be stored in a table is wider than TOAST_TUPLE_THRESHOLD bytes (normally 2 kB). The TOAST code will compress and/or move field values out-of-line until the row value is shorter than TOAST_TUPLE_TARGET bytes (also normally 2 kB, adjustable) or no more gains can be had. During an UPDATE operation, values of unchanged fields are normally preserved as-is; so an UPDATE of a row with out-of-line values incurs no TOAST costs if none of the out-of-line values change.

The TOAST management code recognizes four different strategies for storing TOAST-able columns on disk:

- **PLAIN** prevents either compression or out-of-line storage; furthermore it disables use of single-byte headers for varlena types. This is the only possible strategy for columns of non-TOAST-able data types.
EXTENDED allows both compression and out-of-line storage. This is the default for most TOASTable data types. Compression will be attempted first, then out-of-line storage if the row is still too big.

EXTERNAL allows out-of-line storage but not compression. Use of EXTERNAL will make substring operations on wide text and bytea columns faster (at the penalty of increased storage space) because these operations are optimized to fetch only the required parts of the out-of-line value when it is not compressed.

MAIN allows compression but not out-of-line storage. (Actually, out-of-line storage will still be performed for such columns, but only as a last resort when there is no other way to make the row small enough to fit on a page.)

Each TOAST-able data type specifies a default strategy for columns of that data type, but the strategy for a given table column can be altered with ALTER TABLE ... SET STORAGE.

TOAST_TUPLE_TARGET can be adjusted for each table using ALTER TABLE ... SET (toast_tuple_target = N)

This scheme has a number of advantages compared to a more straightforward approach such as allowing row values to span pages. Assuming that queries are usually qualified by comparisons against relatively small key values, most of the work of the executor will be done using the main row entry. The big values of TOASTed attributes will only be pulled out (if selected at all) at the time the result set is sent to the client. Thus, the main table is much smaller and more of its rows fit in the shared buffer cache than would be the case without any out-of-line storage. Sorts will shrink also, and sorts will more often be done entirely in memory. A little test showed that a table containing typical HTML pages and their URLs was stored in about half of the raw data size including the TOAST table, and that the main table contained only about 10% of the entire data (the URLs and some small HTML pages). There was no run time difference compared to an un-TOASTed comparison table, in which all the HTML pages were cut down to 7 kB to fit.

68.2.2. Out-of-line, in-memory TOAST storage

TOAST pointers can point to data that is not on disk, but is elsewhere in the memory of the current server process. Such pointers obviously cannot be long-lived, but they are nonetheless useful. There are currently two sub-cases: pointers to indirect data and pointers to expanded data.

Indirect TOAST pointers simply point at a non-indirect varlena value stored somewhere in memory. This case was originally created merely as a proof of concept, but it is currently used during logical decoding to avoid possibly having to create physical tuples exceeding 1 GB (as pulling all out-of-line field values into the tuple might do). The case is of limited use since the creator of the pointer datum is entirely responsible that the referenced data survives for as long as the pointer could exist, and there is no infrastructure to help with this.

Expanded TOAST pointers are useful for complex data types whose on-disk representation is not especially suited for computational purposes. As an example, the standard varlena representation of a PostgreSQL array includes dimensionality information, a nulls bitmap if there are any null elements, then the values of all the elements in order. When the element type itself is variable-length, the only way to find the Nth element is to scan through all the preceding elements. This representation is appropriate for on-disk storage because of its compactness, but for computations with the array it's much nicer to have an "expanded" or "deconstructed" representation in which all the element starting locations have been identified. The TOAST pointer mechanism supports this need by allowing a pass-by-reference Datum to point to either a standard varlena value (the on-disk representation) or a TOAST pointer that points to an expanded representation somewhere in memory. The details of this expanded representation are up to the data type, though it must have a standard header and meet the other API requirements given in src/include/utils/expandeddatum.h. C-level functions working with the data type can choose to handle either representation. Functions that do not know about the expanded representation, but simply apply PG_DETOST_DATUM to their inputs, will automatically receive the traditional varlena representation; so support for an expanded representation can be introduced incrementally, one function at a time.
Database Physical Storage

TOAST pointers to expanded values are further broken down into read-write and read-only pointers. The pointed-to representation is the same either way, but a function that receives a read-write pointer is allowed to modify the referenced value in-place, whereas one that receives a read-only pointer must not; it must first create a copy if it wants to make a modified version of the value. This distinction and some associated conventions make it possible to avoid unnecessary copying of expanded values during query execution.

For all types of in-memory TOAST pointer, the TOAST management code ensures that no such pointer datum can accidentally get stored on disk. In-memory TOAST pointers are automatically expanded to normal in-line varlena values before storage — and then possibly converted to on-disk TOAST pointers, if the containing tuple would otherwise be too big.

68.3. Free Space Map

Each heap and index relation, except for hash indexes, has a Free Space Map (FSM) to keep track of available space in the relation. It's stored alongside the main relation data in a separate relation fork, named after the filenode number of the relation, plus a _fsm suffix. For example, if the filenode of a relation is 12345, the FSM is stored in a file called 12345_fsm, in the same directory as the main relation file.

The Free Space Map is organized as a tree of FSM pages. The bottom level FSM pages store the free space available on each heap (or index) page, using one byte to represent each such page. The upper levels aggregate information from the lower levels.

Within each FSM page is a binary tree, stored in an array with one byte per node. Each leaf node represents a heap page, or a lower level FSM page. In each non-leaf node, the higher of its children's values is stored. The maximum value in the leaf nodes is therefore stored at the root.

See src/backend/storage/freespace/README for more details on how the FSM is structured, and how it's updated and searched. The pg_freespacemap module can be used to examine the information stored in free space maps.

68.4. Visibility Map

Each heap relation has a Visibility Map (VM) to keep track of which pages contain only tuples that are known to be visible to all active transactions; it also keeps track of which pages contain only frozen tuples. It's stored alongside the main relation data in a separate relation fork, named after the filenode number of the relation, plus a _vm suffix. For example, if the filenode of a relation is 12345, the VM is stored in a file called 12345_vm, in the same directory as the main relation file. Note that indexes do not have VMs.

The visibility map stores two bits per heap page. The first bit, if set, indicates that the page is all-visible, or in other words that the page does not contain any tuples that need to be vacuumed. This information can also be used by index-only scans to answer queries using only the index tuple. The second bit, if set, means that all tuples on the page have been frozen. That means that even an anti-wraparound vacuum need not revisit the page.

The map is conservative in the sense that we make sure that whenever a bit is set, we know the condition is true, but if a bit is not set, it might or might not be true. Visibility map bits are only set by vacuum, but are cleared by any data-modifying operations on a page.

The pg_visibility module can be used to examine the information stored in the visibility map.

68.5. The Initialization Fork

Each unlogged table, and each index on an unlogged table, has an initialization fork. The initialization fork is an empty table or index of the appropriate type. When an unlogged table must be reset to empty due to a crash, the initialization fork is copied over the main fork, and any other forks are erased (they will be recreated automatically as needed).
68.6. Database Page Layout

This section provides an overview of the page format used within PostgreSQL tables and indexes. Sequences and TOAST tables are formatted just like a regular table.

In the following explanation, a byte is assumed to contain 8 bits. In addition, the term item refers to an individual data value that is stored on a page. In a table, an item is a row; in an index, an item is an index entry.

Every table and index is stored as an array of pages of a fixed size (usually 8 kB, although a different page size can be selected when compiling the server). In a table, all the pages are logically equivalent, so a particular item (row) can be stored in any page. In indexes, the first page is generally reserved as a metapage holding control information, and there can be different types of pages within the index, depending on the index access method.

Table 68.2 shows the overall layout of a page. There are five parts to each page.

### Table 68.2. Overall Page Layout

<table>
<thead>
<tr>
<th>Item</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PageHeaderData</td>
<td>24 bytes long. Contains general information about the page, including free space pointers.</td>
</tr>
<tr>
<td>ItemIdData</td>
<td>Array of (offset,length) pairs pointing to the actual items. 4 bytes per item.</td>
</tr>
<tr>
<td>Free space</td>
<td>The unallocated space. New item pointers are allocated from the start of this area, new items from the end.</td>
</tr>
<tr>
<td>Items</td>
<td>The actual items themselves.</td>
</tr>
<tr>
<td>Special space</td>
<td>Index access method specific data. Different methods store different data. Empty in ordinary tables.</td>
</tr>
</tbody>
</table>

The first 24 bytes of each page consists of a page header (PageHeaderData). Its format is detailed in Table 68.3. The first field tracks the most recent WAL entry related to this page. The second field contains the page checksum if data checksums are enabled. Next is a 2-byte field containing flag bits. This is followed by three 2-byte integer fields (pd_lower, pd_upper, and pd_special). These contain byte offsets from the page start to the start of unallocated space, to the end of unallocated space, and to the start of the special space. The next 2 bytes of the page header, pd_pagesize_version, store both the page size and a version indicator. Beginning with PostgreSQL 8.3 the version number is 4; PostgreSQL 8.1 and 8.2 used version number 3; PostgreSQL 8.0 used version number 2; PostgreSQL 7.3 and 7.4 used version number 1; prior releases used version number 0. (The basic page layout and header format has not changed in most of these versions, but the layout of heap row headers has.) The page size is basically only present as a cross-check; there is no support for having more than one page size in an installation. The last field is a hint that shows whether pruning the page is likely to be profitable: it tracks the oldest un-pruned XMAX on the page.

### Table 68.3. PageHeaderData Layout

<table>
<thead>
<tr>
<th>Field</th>
<th>Type</th>
<th>Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pd_lsn</td>
<td>PageXLogRecPtr</td>
<td>8 bytes</td>
<td>LSN: next byte after last byte of WAL record for last change to this page</td>
</tr>
<tr>
<td>pd_checksum</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Page checksum</td>
</tr>
</tbody>
</table>

1 Actually, index access methods need not use this page format. All the existing index methods do use this basic format, but the data kept on index metapages usually doesn't follow the item layout rules.
Database Physical Storage

<table>
<thead>
<tr>
<th>Field</th>
<th>Type</th>
<th>Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pd_flags</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Flag bits</td>
</tr>
<tr>
<td>pd_lower</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to start of free space</td>
</tr>
<tr>
<td>pd_upper</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to end of free space</td>
</tr>
<tr>
<td>pd_special</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to start of special space</td>
</tr>
<tr>
<td>pd_pagesize_version</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Page size and layout version number information</td>
</tr>
<tr>
<td>pd_prune_xid</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>Oldest unpruned XMAX on page, or zero if none</td>
</tr>
</tbody>
</table>

All the details can be found in src/include/storage/bufpage.h.

Following the page header are item identifiers (ItemIdData), each requiring four bytes. An item identifier contains a byte-offset to the start of an item, its length in bytes, and a few attribute bits which affect its interpretation. New item identifiers are allocated as needed from the beginning of the unallocated space. The number of item identifiers present can be determined by looking at pd_lower, which is increased to allocate a new identifier. Because an item identifier is never moved until it is freed, its index can be used on a long-term basis to reference an item, even when the item itself is moved around on the page to compact free space. In fact, every pointer to an item (ItemPointer, also known as CTID) created by PostgreSQL consists of a page number and the index of an item identifier.

The items themselves are stored in space allocated backwards from the end of unallocated space. The exact structure varies depending on what the table is to contain. Tables and sequences both use a structure named HeapTupleHeaderData, described below.

The final section is the “special section” which can contain anything the access method wishes to store. For example, b-tree indexes store links to the page's left and right siblings, as well as some other data relevant to the index structure. Ordinary tables do not use a special section at all (indicated by setting pd_special to equal the page size).

### 68.6.1. Table Row Layout

All table rows are structured in the same way. There is a fixed-size header (occupying 23 bytes on most machines), followed by an optional null bitmap, an optional object ID field, and the user data. The header is detailed in Table 68.4. The actual user data (columns of the row) begins at the offset indicated by t_hoff, which must always be a multiple of the MAXALIGN distance for the platform. The null bitmap is only present if the HEAP_HASNULL bit is set in t_infomask. If it is present it begins just after the fixed header and occupies enough bytes to have one bit per data column (that is, t_natts bits altogether). In this list of bits, a 1 bit indicates not-null, a 0 bit is a null. When the bitmap is not present, all columns are assumed not-null. The object ID is only present if the HEAP_HASOID bit is set in t_infomask. If present, it appears just before the t_hoff boundary. Any padding needed to make t_hoff a MAXALIGN multiple will appear between the null bitmap and the object ID. (This in turn ensures that the object ID is suitably aligned.)

#### Table 68.4. HeapTupleHeaderData Layout

<table>
<thead>
<tr>
<th>Field</th>
<th>Type</th>
<th>Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>t_xmin</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>insert XID stamp</td>
</tr>
<tr>
<td>t_xmax</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>delete XID stamp</td>
</tr>
<tr>
<td>Field</td>
<td>Type</td>
<td>Length</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------------</td>
<td>--------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>t_cid</td>
<td>CommandId</td>
<td>4 bytes</td>
<td>insert and/or delete CID stamp (overlays with t_xvac)</td>
</tr>
<tr>
<td>t_xvac</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>XID for VACUUM operation moving a row version</td>
</tr>
<tr>
<td>t_ctid</td>
<td>ItemPointerData</td>
<td>6 bytes</td>
<td>current TID of this or newer row version</td>
</tr>
<tr>
<td>t_infomask2</td>
<td>uint16</td>
<td>2 bytes</td>
<td>number of attributes, plus various flag bits</td>
</tr>
<tr>
<td>t_infomask</td>
<td>uint16</td>
<td>2 bytes</td>
<td>various flag bits</td>
</tr>
<tr>
<td>t_hoff</td>
<td>uint8</td>
<td>1 byte</td>
<td>offset to user data</td>
</tr>
</tbody>
</table>

All the details can be found in `src/include/access/htup_details.h`.

Interpreting the actual data can only be done with information obtained from other tables, mostly `pg_attribute`. The key values needed to identify field locations are `attlen` and `attalign`. There is no way to directly get a particular attribute, except when there are only fixed width fields and no null values. All this trickery is wrapped up in the functions `heap_getattr`, `fastgetattr` and `heap_getsysattr`.

To read the data you need to examine each attribute in turn. First check whether the field is NULL according to the null bitmap. If it is, go to the next. Then make sure you have the right alignment. If the field is a fixed width field, then all the bytes are simply placed. If it's a variable length field (attlen = -1) then it's a bit more complicated. All variable-length data types share the common header structure `struct varlena`, which includes the total length of the stored value and some flag bits. Depending on the flags, the data can be either inline or in a TOAST table; it might be compressed, too (see Section 68.2).
Chapter 69. System Catalog

Declarations and Initial Contents

PostgreSQL uses many different system catalogs to keep track of the existence and properties of database objects, such as tables and functions. Physically there is no difference between a system catalog and a plain user table, but the backend C code knows the structure and properties of each catalog, and can manipulate it directly at a low level. Thus, for example, it is inadvisable to attempt to alter the structure of a catalog on-the-fly; that would break assumptions built into the C code about how rows of the catalog are laid out. But the structure of the catalogs can change between major versions.

The structures of the catalogs are declared in specially formatted C header files in the `src/include/catalog/` directory of the source tree. In particular, for each catalog there is a header file named after the catalog (e.g., `pg_class.h` for `pg_class`), which defines the set of columns the catalog has, as well as some other basic properties such as its OID. Other critical files defining the catalog structure include `indexing.h`, which defines the indexes present on all the system catalogs, and `toasting.h`, which defines TOAST tables for catalogs that need one.

Many of the catalogs have initial data that must be loaded into them during the “bootstrap” phase of `initdb`, to bring the system up to a point where it is capable of executing SQL commands. (For example, `pg_class.h` must contain an entry for itself, as well as one for each other system catalog and index.) This initial data is kept in editable form in data files that are also stored in the `src/include/catalog/` directory. For example, `pg_proc.dat` describes all the initial rows that must be inserted into the `pg_proc` catalog.

To create the catalog files and load this initial data into them, a backend running in bootstrap mode reads a BKI (Backend Interface) file containing commands and initial data. The `postgres.bki` file used in this mode is prepared from the aforementioned header and data files, while building a PostgreSQL distribution, by a Perl script named `genbki.pl`. Although it's specific to a particular PostgreSQL release, `postgres.bki` is platform-independent and is installed in the `share` subdirectory of the installation tree.

`genbki.pl` also produces a derived header file for each catalog, for example `pg_class_d.h` for the `pg_class` catalog. This file contains automatically-generated macro definitions, and may contain other macros, enum declarations, and so on that can be useful for client C code that reads a particular catalog.

Most Postgres developers don’t need to be directly concerned with the BKI file, but almost any non-trivial feature addition in the backend will require modifying the catalog header files and/or initial data files. The rest of this chapter gives some information about that, and for completeness describes the BKI file format.

### 69.1. System Catalog Declaration Rules

The key part of a catalog header file is a C structure definition describing the layout of each row of the catalog. This begins with a `CATALOG` macro, which so far as the C compiler is concerned is just shorthand for `typedef struct FormData_catalogname`. Each field in the struct gives rise to a catalog column. Fields can be annotated using the BKI property macros described in `genbki.h`, for example to define a default value for a field or mark it as nullable or not nullable. The `CATALOG` line can also be annotated, with some other BKI property macros described in `genbki.h`, to define other properties of the catalog as a whole, such as whether it has OIDs (by default, it does).

The system catalog cache code (and most catalog-munging code in general) assumes that the fixed-length portions of all system catalog tuples are in fact present, because it maps this C struct declaration onto them. Thus, all variable-length fields and nullable fields must be placed at the end, and they cannot be accessed as struct fields. For example, if you tried to set `pg_type.typrelid` to be `NULL`, it would fail when some piece of code tried to reference `typetup->typrelid` (or worse,
typetup->typelem, because that follows typrelid). This would result in random errors or even segmentation violations.

As a partial guard against this type of error, variable-length or nullable fields should not be made directly visible to the C compiler. This is accomplished by wrapping them in #ifdef CATALOG_VARLEN ... #endif (where CATALOG_VARLEN is a symbol that is never defined). This prevents C code from carelessly trying to access fields that might not be there or might be at some other offset. As an independent guard against creating incorrect rows, we require all columns that should be non-nullable to be marked so in pg_attribute. The bootstrap code will automatically mark catalog columns as NOT NULL if they are fixed-width and are not preceded by any nullable column. Where this rule is inadequate, you can force correct marking by using BKI_FORCE_NOT_NULL and BKI_FORCE_NULL annotations as needed. But note that NOT NULL constraints are only enforced in the executor, not against tuples that are generated by random C code, so care is still needed when manually creating or updating catalog rows.

Frontend code should not include any pg_xxx.h catalog header file, as these files may contain C code that won't compile outside the backend. (Typically, that happens because these files also contain declarations for functions in src/backend/catalog/files.) Instead, frontend code may include the corresponding generated pg_xxx_d.h header, which will contain OID #defines and any other data that might be of use on the client side. If you want macros or other code in a catalog header to be visible to frontend code, write #ifdef EXPOSE_TO_CLIENT_CODE ... #endif around that section to instruct genbki.pl to copy that section to the pg_xxx_d.h header.

A few of the catalogs are so fundamental that they can't even be created by the BKI create command that's used for most catalogs, because that command needs to write information into these catalogs to describe the new catalog. These are called bootstrap catalogs, and defining one takes a lot of extra work: you have to manually prepare appropriate entries for them in the pre-loaded contents of pg_class and pg_type, and those entries will need to be updated for subsequent changes to the catalog's structure. (Bootstrap catalogs also need pre-loaded entries in pg_attribute, but fortunately genbki.pl handles that chore nowadays.) Avoid making new catalogs be bootstrap catalogs if at all possible.

### 69.2. System Catalog Initial Data

Each catalog that has any manually-created initial data (some do not) has a corresponding .dat file that contains its initial data in an editable format.

#### 69.2.1. Data File Format

Each .dat file contains Perl data structure literals that are simply eval'd to produce an in-memory data structure consisting of an array of hash references, one per catalog row. A slightly modified excerpt from pg_database.dat will demonstrate the key features:

```
[
  # A comment could appear here.
  { oid => '1', oid_symbol => 'TemplateDbOid',
    descr => 'database\'s default template',
    datname => 'template1', datdba => 'PGUID',
    encoding => 'ENCODING',
    datcollate => 'LC_COLLATE',
    datctype => 'LC_CTYPE',
    datistemplate => 't',
    datallowconn => 't',
    datconnlimit => '-1',
    datlastsysoid => '0',
    datfrozenxid => '0',
    datminmxid => '1',
    dattablespace => '1663',
    datacl => '_null_' },
```
Points to note:

- The overall file layout is: open square bracket, one or more sets of curly braces each of which represents a catalog row, close square bracket. Write a comma after each closing curly brace.

- Within each catalog row, write comma-separated key => value pairs. The allowed keys are the names of the catalog's columns, plus the metadata keys oid, oid_symbol, and descr. (The use of oid and oid_symbol is described in Section 69.2.2 below. descr supplies a description string for the object, which will be inserted into pg_description or pg_shdescription as appropriate.) While the metadata keys are optional, the catalog's defined columns must all be provided, except when the catalog's .h file specifies a default value for the column.

- All values must be single-quoted. Escape single quotes used within a value with a backslash. Backslashes meant as data can, but need not, be doubled; this follows Perl's rules for simple quoted literals. Note that backslashes appearing as data will be treated as escapes by the bootstrap scanner, according to the same rules as for escape string constants (see Section 4.1.2.2); for example \t converts to a tab character. If you actually want a backslash in the final value, you will need to write four of them: Perl strips two, leaving \\ for the bootstrap scanner to see.

- Null values are represented by _null_. (Note that there is no way to create a value that is just that string.)

- Comments are preceded by #, and must be on their own lines.

- To aid readability, field values that are OIDs of other catalog entries can be represented by names rather than numeric OIDs. This is described in Section 69.2.3 below.

- Since hashes are unordered data structures, field order and line layout aren't semantically significant. However, to maintain a consistent appearance, we set a few rules that are applied by the formatting script reformat_dat_file.pl:
  - Within each pair of curly braces, the metadata fields oid, oid_symbol, and descr (if present) come first, in that order, then the catalog's own fields appear in their defined order.
  - Newlines are inserted between fields as needed to limit line length to 80 characters, if possible. A newline is also inserted between the metadata fields and the regular fields.
  - If the catalog's .h file specifies a default value for a column, and a data entry has that same value, reformat_dat_file.pl will omit it from the data file. This keeps the data representation compact.

- reformat_dat_file.pl preserves blank lines and comment lines as-is. It's recommended to run reformat_dat_file.pl before submitting catalog data patches. For convenience, you can simply change to src/include/catalog/ and run make refor-
mat-dat-files.

- If you want to add a new method of making the data representation smaller, you must implement it in reformat_dat_file.pl and also teach Catalog::ParseData() how to expand the data back into the full representation.

### 69.2.2. OID Assignment

A catalog row appearing in the initial data can be given a manually-assigned OID by writing an oid => nnnn metadata field. Furthermore, if an OID is assigned, a C macro for that OID can be created by writing an oid_symbol => name metadata field.

Pre-loaded catalog rows must have preassigned OIDs if there are OID references to them in other pre-loaded rows. A preassigned OID is also needed if the row's OID must be referenced from C code. If neither case applies, the oid metadata field can be omitted, in which case the bootstrap code assigns an OID automatically, or leaves it zero in a catalog that has no OIDs. In practice we usually preassign
OIDs for all or none of the pre-loaded rows in a given catalog, even if only some of them are actually cross-referenced.

Writing the actual numeric value of any OID in C code is considered very bad form; always use a macro, instead. Direct references to \texttt{pg\_proc} OIDs are common enough that there's a special mechanism to create the necessary macros automatically; see \texttt{src/backend/utils/Gen_fmgrtab.pl}. Similarly — but, for historical reasons, not done the same way — there's an automatic method for creating macros for \texttt{pg\_type} OIDs. \texttt{oid\_symbol} entries are therefore not necessary in those two catalogs. Likewise, macros for the \texttt{pg\_class} OIDs of system catalogs and indexes are set up automatically. For all other system catalogs, you have to manually specify any macros you need via \texttt{oid\_symbol} entries.

To find an available OID for a new pre-loaded row, run the script \texttt{src/include/catalog/unused\_oids}. It prints inclusive ranges of unused OIDs (e.g., the output line “45-900” means OIDs 45 through 900 have not been allocated yet). Currently, OIDs 1-9999 are reserved for manual assignment; the \texttt{unused\_oids} script simply looks through the catalog headers and .dat files to see which ones do not appear. You can also use the \texttt{duplicate\_oids} script to check for mistakes. (\texttt{genbki.pl} will also detect duplicate OIDs at compile time.)

The OID counter starts at 10000 at the beginning of a bootstrap run. If a catalog row is in a table that requires OIDs, but no OID was preassigned by an \texttt{oid} field, then it will receive an OID of 10000 or above.

### 69.2.3. OID Reference Lookup

Cross-references from one initial catalog row to another can be written by just writing the preassigned OID of the referenced row. But that's error-prone and hard to understand, so for frequently-referenced catalogs, \texttt{genbki.pl} provides mechanisms to write symbolic references instead. Currently this is possible for references to access methods, functions, operators, opclasses, opfamilies, and types. The rules are as follows:

- Use of symbolic references is enabled in a particular catalog column by attaching \texttt{BKI\_LOOKUP(lookuprule)} to the column's definition, where \texttt{lookuprule} is \texttt{pg\_am}, \texttt{pg\_proc}, \texttt{pg\_operator}, \texttt{pg\_opclass}, \texttt{pg\_opfamily}, or \texttt{pg\_type}. \texttt{BKI\_LOOKUP} can be attached to columns of type \texttt{Oid}, \texttt{regproc}, \texttt{oidvector}, or \texttt{Oid[]} in the latter two cases it implies performing a lookup on each element of the array.

- In such a column, all entries must use the symbolic format except when writing \texttt{0} for InvalidOid. (If the column is declared \texttt{regproc}, you can optionally write \texttt{-} instead of \texttt{0}.) \texttt{genbki.pl} will warn about unrecognized names.

- Access methods are just represented by their names, as are types. Type names must match the referenced \texttt{pg\_type} entry's \texttt{typname}; you do not get to use any aliases such as \texttt{integer} for \texttt{int4}.

- A function can be represented by its \texttt{proname}, if that is unique among the \texttt{pg\_proc.dat} entries (this works like \texttt{regproc} input). Otherwise, write it as \texttt{proname(argtypename,argtypename,...)}, like \texttt{regprocedure}. The argument type names must be spelled exactly as they are in the \texttt{pg\_proc.dat} entry's \texttt{proargtypes} field. Do not insert any spaces.

- Operators are represented by \texttt{oprname(lefttype,righttype)}, writing the type names exactly as they appear in the \texttt{pg\_operator.dat} entry's \texttt{oprleft} and \texttt{oprright} fields. (Write \texttt{0} for the omitted operand of a unary operator.)

- The names of opclasses and opfamilies are only unique within an access method, so they are represented by \texttt{access\_method\_name/object\_name}.

- In none of these cases is there any provision for schema-qualification; all objects created during bootstrap are expected to be in the \texttt{pg\_catalog} schema.
genbki.pl resolves all symbolic references while it runs, and puts simple numeric OIDs into the emitted BKI file. There is therefore no need for the bootstrap backend to deal with symbolic references.

## 69.2.4. Recipes for Editing Data Files

Here are some suggestions about the easiest ways to perform common tasks when updating catalog data files.

### Add a new column with a default to a catalog:
Add the column to the header file with a `BKI_DEFAULT(value)` annotation. The data file need only be adjusted by adding the field in existing rows where a non-default value is needed.

### Add a default value to an existing column that doesn’t have one:
Add a `BKI_DEFAULT` annotation to the header file, then run `make reformat-dat-files` to remove now-redundant field entries.

### Remove a column, whether it has a default or not:
Remove the column from the header, then run `make reformat-dat-files` to remove now-useless field entries.

### Change or remove an existing default value:
You cannot simply change the header file, since that will cause the current data to be interpreted incorrectly. First run `make expand-dat-files` to rewrite the data files with all default values inserted explicitly, then change or remove the `BKI_DEFAULT` annotation, then run `make reformat-dat-files` to remove superfluous fields again.

### Ad-hoc bulk editing:
`reformat_dat_file.pl` can be adapted to perform many kinds of bulk changes. Look for its block comments showing where one-off code can be inserted. In the following example, we are going to consolidate two boolean fields in `pg_proc` into a char field:

1. Add the new column, with a default, to `pg_proc.h`:

   ```c
   +    /* see PROKIND_ categories below */
   +    char        prokind BKI_DEFAULT(f);
   ```

2. Create a new script based on `reformat_dat_file.pl` to insert appropriate values on-the-fly:

   ```perl
   -           # At this point we have the full row in memory as a hash
   -           # and can do any operations we want. As written, it only
   -           # removes default values, but this script can be adapted to
   -           # do one-off bulk-editing.
   +           # One-off change to migrate to prokind
   +           # Default has already been filled in by now, so change to other
   +           # values as appropriate
   +           if ($values{proisagg} eq 't')
   +               $values{prokind} = 'a';
   +           }
   +           elsif ($values{proiswindow} eq 't')
   +               $values{prokind} = 'w';
   +           }
   ```

3. Run the new script:
$ cd src/include/catalog
$ perl rewrite_dat_with_prokind.pl pg_proc.dat

At this point pg_proc.dat has all three columns, prokind, proisagg, and proiswindow, though they will appear only in rows where they have non-default values.

4. Remove the old columns from pg_proc.h:

- /* is it an aggregate? */
- bool proisagg BKI_DEFAULT(f);
- /* is it a window function? */
- bool proiswindow BKI_DEFAULT(f);

5. Finally, run make reformat-dat-files to remove the useless old entries from pg_proc.dat.

For further examples of scripts used for bulk editing, see convert_oid2name.pl and remove_pg_type_oid_symbols.pl attached to this message: https://www.postgresql.org/message-id/CAJV6QJGVX8gXnPt+Xa=DxR7kFyprQ1tNcCT5D0O3ShfnM6jehA@mail.gmail.com

69.3. BKI File Format

This section describes how the PostgreSQL backend interprets BKI files. This description will be easier to understand if the postgres.bki file is at hand as an example.

BKI input consists of a sequence of commands. Commands are made up of a number of tokens, depending on the syntax of the command. Tokens are usually separated by whitespace, but need not be if there is no ambiguity. There is no special command separator; the next token that syntactically cannot belong to the preceding command starts a new one. (Usually you would put a new command on a new line, for clarity.) Tokens can be certain key words, special characters (parentheses, commas, etc.), numbers, or double-quoted strings. Everything is case sensitive.

Lines starting with # are ignored.

69.4. BKI Commands

create tablename tableoid [bootstrap] [shared_relation] [without_oids] [rowtype_oid oid](name1=type1[FORCE NOT NULL|FORCE NULL][, name2=type2 [FORCE NOT NULL|FORCE NULL],...))

Create a table named tablename, and having the OID tableoid, with the columns given in parentheses.

The following column types are supported directly by bootstrap.c: bool, bytea, char (1 byte), name, int2, int4, regproc, regclass, regtype, text, oid, tid, xid, cid, int2vector, oidvector, _int4 (array), _text (array), _oid (array), _char (array), _aclitem (array). Although it is possible to create tables containing columns of other types, this cannot be done until after pg_type has been created and filled with appropriate entries. (That effectively means that only these column types can be used in bootstrap catalogs, but non-bootstrap catalogs can contain any built-in type.)

When bootstrap is specified, the table will only be created on disk; nothing is entered into pg_class, pg_attribute, etc, for it. Thus the table will not be accessible by ordinary SQL operations until such entries are made the hard way (with insert commands). This option is used for creating pg_class etc themselves.
The table is created as shared if shared_relation is specified. It will have OIDs unless without_oids is specified. The table's row type OID (pg_type OID) can optionally be specified via the rowtype_oid clause; if not specified, an OID is automatically generated for it. (The rowtype_oid clause is useless if bootstrap is specified, but it can be provided anyway for documentation.)

open tablename

Open the table named tablename for insertion of data. Any currently open table is closed.

close tablename

Close the open table. The name of the table must be given as a cross-check.

insert [OID = oid_value] ( value1 value2 ...)

Insert a new row into the open table using value1, value2, etc., for its column values and oid_value for its OID. If oid_value is zero (0) or the clause is omitted, and the table has OIDs, then the next available OID is assigned.

NULL values can be specified using the special key word _null_. Values that do not look like identifiers or digit strings must be double quoted.

declare [unique] index indexname indexoid on tablename using amname ( opclass1 name1[, ...])

Create an index named indexname, having OID indexoid, on the table named tablename, using the amname access method. The fields to index are called name1, name2 etc., and the operator classes to use are opclass1, opclass2 etc., respectively. The index file is created and appropriate catalog entries are made for it, but the index contents are not initialized by this command.

declare toast toasttableoid toastindexoid on tablename

Create a TOAST table for the table named tablename. The TOAST table is assigned OID toasttableoid and its index is assigned OID toastindexoid. As with declare index, filling of the index is postponed.

build indices

Fill in the indices that have previously been declared.

69.5. Structure of the Bootstrap BKI File

The open command cannot be used until the tables it uses exist and have entries for the table that is to be opened. (These minimum tables are pg_class, pg_attribute, pg_proc, and pg_type.) To allow those tables themselves to be filled, create with the bootstrap option implicitly opens the created table for data insertion.

Also, the declare index and declare toast commands cannot be used until the system catalogs they need have been created and filled in.

Thus, the structure of the postgres.bki file has to be:

1. create bootstrap one of the critical tables

2. insert data describing at least the critical tables

3. close

4. Repeat for the other critical tables.
5. create (without bootstrap) a noncritical table
6. open
7. insert desired data
8. close
9. Repeat for the other noncritical tables.
10. Define indexes and toast tables.
11. build indices

There are doubtless other, undocumented ordering dependencies.

69.6. BKI Example

The following sequence of commands will create the table test_table with OID 420, having two columns cola and colb of type int4 and text, respectively, and insert two rows into the table:

```sql
create test_table 420 (cola = int4, colb = text)
open test_table
insert OID=421 ( 1 "value1" )
insert OID=422 ( 2 _null_ )
close test_table
```
Chapter 70. How the Planner Uses Statistics

This chapter builds on the material covered in Section 14.1 and Section 14.2 to show some additional details about how the planner uses the system statistics to estimate the number of rows each part of a query might return. This is a significant part of the planning process, providing much of the raw material for cost calculation.

The intent of this chapter is not to document the code in detail, but to present an overview of how it works. This will perhaps ease the learning curve for someone who subsequently wishes to read the code.

70.1. Row Estimation Examples

The examples shown below use tables in the PostgreSQL regression test database. The outputs shown are taken from version 8.3. The behavior of earlier (or later) versions might vary. Note also that since \texttt{ANALYZE} uses random sampling while producing statistics, the results will change slightly after any \texttt{ANALYZE}.

Let's start with a very simple query:

\begin{verbatim}
EXPLAIN SELECT * FROM tenk1;
\end{verbatim}

\begin{verbatim}
QUERY PLAN
-------------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..458.00 rows=10000 width=244)
How the planner determines the cardinality of \texttt{tenk1} is covered in Section 14.2, but is repeated here for completeness. The number of pages and rows is looked up in \texttt{pg_class}:

\begin{verbatim}
SELECT relpages, reltuples FROM pg_class WHERE relname = 'tenk1';
\end{verbatim}

\begin{verbatim}
relpages | reltuples
----------+-----------
358 |     10000
\end{verbatim}

These numbers are current as of the last \texttt{VACUUM} or \texttt{ANALYZE} on the table. The planner then fetches the actual current number of pages in the table (this is a cheap operation, not requiring a table scan). If that is different from \texttt{repages} then \texttt{reltuples} is scaled accordingly to arrive at a current number-of-rows estimate. In the example above, the value of \texttt{repages} is up-to-date so the rows estimate is the same as \texttt{reltuples}.

Let's move on to an example with a range condition in its \texttt{WHERE} clause:

\begin{verbatim}
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 1000;
\end{verbatim}

\begin{verbatim}
QUERY PLAN
--------------------------------------------------------------------------------
Bitmap Heap Scan on tenk1  (cost=24.06..394.64 rows=1007 width=244)
  Recheck Cond: (unique1 < 1000)
  ->  Bitmap Index Scan on tenk1_unique1  (cost=0.00..23.80 rows=1007 width=0)
\end{verbatim}
How the Planner Uses Statistics

Index Cond: (unique1 < 1000)

The planner examines the WHERE clause condition and looks up the selectivity function for the operator < in pg_operator. This is held in the column oprrest, and the entry in this case is scalarltsel. The scalarltsel function retrieves the histogram for unique1 from pg_statistic. For manual queries it is more convenient to look in the simpler pg_stats view:

```
SELECT histogram_bounds FROM pg_stats
WHERE tablename='tenk1' AND attname='unique1';
```

<table>
<thead>
<tr>
<th>histogram_bounds</th>
</tr>
</thead>
<tbody>
<tr>
<td>(0,993,1997,3050,4040,5036,5957,7057,8029,9016,9995)</td>
</tr>
</tbody>
</table>

Next the fraction of the histogram occupied by “< 1000” is worked out. This is the selectivity. The histogram divides the range into equal frequency buckets, so all we have to do is locate the bucket that our value is in and count part of it and all of the ones before. The value 1000 is clearly in the second bucket (993-1997). Assuming a linear distribution of values inside each bucket, we can calculate the selectivity as:

```
selectivity = (1 + (1000 - bucket[2].min)/(bucket[2].max - bucket[2].min))/num_buckets
  = (1 + (1000 - 993)/(1997 - 993))/10
  = 0.100697
```

that is, one whole bucket plus a linear fraction of the second, divided by the number of buckets. The estimated number of rows can now be calculated as the product of the selectivity and the cardinality of tenk1:

```
rows = rel_cardinality * selectivity
  = 10000 * 0.100697
  = 1007  (rounding off)
```

Next let's consider an example with an equality condition in its WHERE clause:

```
EXPLAIN SELECT * FROM tenk1 WHERE stringu1 = 'CRAAAA';
```

```
QUERY PLAN
----------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..483.00 rows=30 width=244)
  Filter: (stringu1 = 'CRAAAA':::name)
```

Again the planner examines the WHERE clause condition and looks up the selectivity function for =, which is eqsel. For equality estimation the histogram is not useful; instead the list of most common values (MCVs) is used to determine the selectivity. Let's have a look at the MCVs, with some additional columns that will be useful later:

```
SELECT null_frac, n_distinct, most_common_vals, most_common_freqs
FROM pg_stats
WHERE tablename='tenk1' AND attname='stringu1';
```

<table>
<thead>
<tr>
<th>null_frac</th>
<th>n_distinct</th>
<th>most_common_vals</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>676</td>
<td>(EJAAAA, BBAAAA, CRAAAA, FCAAAA, FEAAAA, GSAAAA, JOAAAA, MCAAAA, NAAAAA, WAAAAA)</td>
</tr>
</tbody>
</table>

2259
How the Planner Uses Statistics

most_common_freqs |
     {0.00333333, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003}

Since CRAAAA appears in the list of MCVs, the selectivity is merely the corresponding entry in the list of most common frequencies (MCFs):

\[ \text{selectivity} = \text{mcf}[3] = 0.003 \]

As before, the estimated number of rows is just the product of this with the cardinality of tenk1:

\[ \text{rows} = 10000 \times 0.003 = 30 \]

Now consider the same query, but with a constant that is not in the MCV list:

EXPLAIN SELECT * FROM tenk1 WHERE stringu1 = 'xxx';

QUERY PLAN

Seq Scan on tenk1  (cost=0.00..483.00 rows=15 width=244)
Filter: (stringu1 = 'xxx'::name)

This is quite a different problem: how to estimate the selectivity when the value is not in the MCV list. The approach is to use the fact that the value is not in the list, combined with the knowledge of the frequencies for all of the MCVs:

\[ \text{selectivity} = \frac{1 - \text{sum(mvf)}}{\text{num_distinct} - \text{num_mcv}} = \frac{1 - (0.00333333 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003)}{676 - 10} = 0.0014559 \]

That is, add up all the frequencies for the MCVs and subtract them from one, then divide by the number of other distinct values. This amounts to assuming that the fraction of the column that is not any of the MCVs is evenly distributed among all the other distinct values. Notice that there are no null values so we don't have to worry about those (otherwise we'd subtract the null fraction from the numerator as well). The estimated number of rows is then calculated as usual:

\[ \text{rows} = 10000 \times 0.0014559 = 15 \text{ (rounding off)} \]

The previous example with unique1 < 1000 was an oversimplification of what scalarltssel really does; now that we have seen an example of the use of MCVs, we can fill in some more detail. The example was correct as far as it went, because since unique1 is a unique column it has no MCVs (obviously, no value is any more common than any other value). For a non-unique column, there will normally be both a histogram and an MCV list, and the histogram does not include the portion of the column population represented by the MCVs. We do things this way because it allows more precise estimation. In this situation scalarltssel directly applies the condition (e.g., "< 1000") to each value of the MCV list, and adds up the frequencies of the MCVs for which the condition is true. This gives an exact estimate of the selectivity within the portion of the table that is MCVs. The histogram is then used in the same way as above to estimate the selectivity in the portion of the table that is not MCVs, and then the two numbers are combined to estimate the overall selectivity. For example, consider
EXPLAIN SELECT * FROM tenk1 WHERE stringu1 < 'IAAAAA';

QUERY PLAN
------------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..483.00 rows=3077 width=244)
  Filter: (stringu1 < 'IAAAAA'::name)

We already saw the MCV information for stringu1, and here is its histogram:

SELECT histogram_bounds FROM pg_stats
WHERE tablename='tenk1' AND attname='stringu1';

<table>
<thead>
<tr>
<th>histogram_bounds</th>
</tr>
</thead>
<tbody>
<tr>
<td>{AAAAAA, CQAAAA, FRAAAA, IBAAAA, KRAAAA, NFAAAA, PSAAAA, SGAAAA, VAAAAA, XLAAAA, ZZAAAA}</td>
</tr>
</tbody>
</table>

Checking the MCV list, we find that the condition stringu1 < 'IAAAAA' is satisfied by the first six entries and not the last four, so the selectivity within the MCV part of the population is

selectivity = sum(relevant mvfs)
= 0.00333333 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003
= 0.01833333

Summing all the MCFs also tells us that the total fraction of the population represented by MCVs is 0.03033333, and therefore the fraction represented by the histogram is 0.96966667 (again, there are no nulls, else we'd have to exclude them here). We can see that the value IAAAAA falls nearly at the end of the third histogram bucket. Using some rather cheesy assumptions about the frequency of different characters, the planner arrives at the estimate 0.298387 for the portion of the histogram population that is less than IAAAAA. We then combine the estimates for the MCV and non-MCV populations:

selectivity = mcv_selectivity + histogram_selectivity * histogram_fraction
= 0.01833333 + 0.298387 * 0.96966667
= 0.307669
rows = 10000 * 0.307669
= 3077  (rounding off)

In this particular example, the correction from the MCV list is fairly small, because the column distribution is actually quite flat (the statistics showing these particular values as being more common than others are mostly due to sampling error). In a more typical case where some values are significantly more common than others, this complicated process gives a useful improvement in accuracy because the selectivity for the most common values is found exactly.

Now let's consider a case with more than one condition in the WHERE clause:

EXPLAIN SELECT * FROM tenk1 WHERE uniquel < 1000 AND stringu1 = 'xxx';

QUERY PLAN
------------------------------------------------------------
Bitmap Heap Scan on tenk1  (cost=23.80..396.91 rows=1 width=244)
  Recheck Cond: (uniquel < 1000)
  Filter: (stringu1 = 'xxx'::name)
The planner assumes that the two conditions are independent, so that the individual selectivities of the clauses can be multiplied together:

\[
\text{selectivity} = \text{selectivity} (\text{unique1} < 1000) \times \text{selectivity} (\text{stringu1} = 'xxx')
\]

\[
= 0.100697 \times 0.0014559 \\
= 0.0001466
\]

\[
\text{rows} = 10000 \times 0.0001466 \\
= 1 \text{ (rounding off)}
\]

Notice that the number of rows estimated to be returned from the bitmap index scan reflects only the condition used with the index; this is important since it affects the cost estimate for the subsequent heap fetches.

Finally we will examine a query that involves a join:

EXPLAIN SELECT * FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 50 AND t1.unique2 = t2.unique2;

The restriction on \text{tenk1}, unique1 < 50, is evaluated before the nested-loop join. This is handled analogously to the previous range example. This time the value 50 falls into the first bucket of the unique1 histogram:

\[
\text{selectivity} = (0 + (50 - \text{bucket}[1].\text{min})/(\text{bucket}[1].\text{max} - \text{bucket}[1].\text{min}))/\text{num_buckets}
\]

\[
= (0 + (50 - 0)/(993 - 0))/10 \\
= 0.005035
\]

\[
\text{rows} = 10000 \times 0.005035 \\
= 50 \text{ (rounding off)}
\]

The restriction for the join is \text{t2.unique2} = \text{t1.unique2}. The operator is just our familiar =, however the selectivity function is obtained from the oprjoin column of \text{pg_operator}, and is eqjoinsel. eqjoinsel looks up the statistical information for both \text{tenk2} and \text{tenk1}:

SELECT tablename, null_frac, n_distinct, most_common_vals FROM pg_stats
WHERE tablename IN ('tenk1', 'tenk2') AND attname='unique2';

<table>
<thead>
<tr>
<th>tablename</th>
<th>null_frac</th>
<th>n_distinct</th>
<th>most_common_vals</th>
</tr>
</thead>
<tbody>
<tr>
<td>tenk1</td>
<td>0</td>
<td>-1</td>
<td></td>
</tr>
<tr>
<td>tenk2</td>
<td>0</td>
<td>-1</td>
<td></td>
</tr>
</tbody>
</table>

In this case there is no MCV information for unique2 because all the values appear to be unique, so we use an algorithm that relies only on the number of distinct values for both relations together with their null fractions:

\[
\text{selectivity} = (1 - \text{null_frac1}) \times (1 - \text{null_frac2}) \times \min\left(\frac{1}{\text{num_distinct1}}, \frac{1}{\text{num_distinct2}}\right)
\]

\[
= (1 - 0) \times (1 - 0) / \max(10000, 10000)
\]

\[
= 0.0001
\]

This is, subtract the null fraction from one for each of the relations, and divide by the maximum of the numbers of distinct values. The number of rows that the join is likely to emit is calculated as the cardinality of the Cartesian product of the two inputs, multiplied by the selectivity:

\[
\text{rows} = (\text{outer_cardinality} \times \text{inner_cardinality}) \times \text{selectivity}
\]

\[
= (50 \times 10000) \times 0.0001
\]

\[
= 50
\]

Had there been MCV lists for the two columns, eqjoinsel would have used direct comparison of the MCV lists to determine the join selectivity within the part of the column populations represented by the MCVs. The estimate for the remainder of the populations follows the same approach shown here.

Notice that we showed inner_cardinality as 10000, that is, the unmodified size of tenk2. It might appear from inspection of the EXPLAIN output that the estimate of join rows comes from 50 * 1, that is, the number of outer rows times the estimated number of rows obtained by each inner index scan on tenk2. But this is not the case: the join relation size is estimated before any particular join plan has been considered. If everything is working well then the two ways of estimating the join size will produce about the same answer, but due to round-off error and other factors they sometimes diverge significantly.

For those interested in further details, estimation of the size of a table (before any WHERE clauses) is done in src/backend/optimizer/util/plancat.c. The generic logic for clause selectivities is in src/backend/optimizer/path/clausesel.c. The operator-specific selectivity functions are mostly found in src/backend/utils/adt/selfuncs.c.

### 70.2. Multivariate Statistics Examples

#### 70.2.1. Functional Dependencies

Multivariate correlation can be demonstrated with a very simple data set — a table with two columns, both containing the same values:

```sql
CREATE TABLE t (a INT, b INT);
INSERT INTO t SELECT i % 100, i % 100 FROM generate_series(1, 10000) s(i);
ANALYZE t;
```

As explained in Section 14.2, the planner can determine cardinality of t using the number of pages and rows obtained from pg_class:
SELECT relpages, reltuples FROM pg_class WHERE relname = 't';

<table>
<thead>
<tr>
<th>relpages</th>
<th>reltuples</th>
</tr>
</thead>
<tbody>
<tr>
<td>45</td>
<td>10000</td>
</tr>
</tbody>
</table>

The data distribution is very simple; there are only 100 distinct values in each column, uniformly distributed.

The following example shows the result of estimating a WHERE condition on the a column:

EXPLAIN (ANALYZE, TIMING OFF) SELECT * FROM t WHERE a = 1;

QUERY PLAN

```
Seq Scan on t  (cost=0.00..170.00 rows=100 width=8) (actual rows=100 loops=1)
  Filter: (a = 1)
  Rows Removed by Filter: 9900
```

The planner examines the condition and determines the selectivity of this clause to be 1%. By comparing this estimate and the actual number of rows, we see that the estimate is very accurate (in fact exact, as the table is very small). Changing the WHERE condition to use the b column, an identical plan is generated. But observe what happens if we apply the same condition on both columns, combining them with AND:

EXPLAIN (ANALYZE, TIMING OFF) SELECT * FROM t WHERE a = 1 AND b = 1;

QUERY PLAN

```
Seq Scan on t  (cost=0.00..195.00 rows=1 width=8) (actual rows=100 loops=1)
  Filter: ((a = 1) AND (b = 1))
  Rows Removed by Filter: 9900
```

The planner estimates the selectivity for each condition individually, arriving at the same 1% estimates as above. Then it assumes that the conditions are independent, and so it multiplies their selectivities, producing a final selectivity estimate of just 0.01%. This is a significant underestimate, as the actual number of rows matching the conditions (100) is two orders of magnitude higher.

This problem can be fixed by creating a statistics object that directs ANALYZE to calculate function-al-dependency multivariate statistics on the two columns:

CREATE STATISTICS stts (dependencies) ON a, b FROM t;
ANALYZE t;

EXPLAIN (ANALYZE, TIMING OFF) SELECT * FROM t WHERE a = 1 AND b = 1;

QUERY PLAN

```
Seq Scan on t  (cost=0.00..195.00 rows=100 width=8) (actual rows=100 loops=1)
  Filter: ((a = 1) AND (b = 1))
  Rows Removed by Filter: 9900
```
70.2.2. Multivariate N-Distinct Counts

A similar problem occurs with estimation of the cardinality of sets of multiple columns, such as the number of groups that would be generated by a \texttt{GROUP BY} clause. When \texttt{GROUP BY} lists a single column, the n-distinct estimate (which is visible as the estimated number of rows returned by the HashAggregate node) is very accurate:

\texttt{EXPLAIN (ANALYZE, TIMING OFF) SELECT COUNT(*) FROM t GROUP BY a;}

\texttt{QUERY PLAN}

\begin{verbatim}
HashAggregate  (cost=195.00..196.00 rows=100 width=12) (actual rows=100 loops=1)
  Group Key: a
    ->  Seq Scan on t  (cost=0.00..145.00 rows=10000 width=4)
         (actual rows=10000 loops=1)
\end{verbatim}

But without multivariate statistics, the estimate for the number of groups in a query with two columns in \texttt{GROUP BY}, as in the following example, is off by an order of magnitude:

\texttt{EXPLAIN (ANALYZE, TIMING OFF) SELECT COUNT(*) FROM t GROUP BY a, b;}

\texttt{QUERY PLAN}

\begin{verbatim}
HashAggregate  (cost=220.00..230.00 rows=1000 width=16) (actual rows=100 loops=1)
  Group Key: a, b
    ->  Seq Scan on t  (cost=0.00..145.00 rows=10000 width=8)
         (actual rows=10000 loops=1)
\end{verbatim}

By redefining the statistics object to include n-distinct counts for the two columns, the estimate is much improved:

\texttt{DROP STATISTICS stts;}
\texttt{CREATE STATISTICS stts (dependencies, ndistinct) ON a, b FROM t;}
\texttt{ANALYZE t;}
\texttt{EXPLAIN (ANALYZE, TIMING OFF) SELECT COUNT(*) FROM t GROUP BY a, b;}

\texttt{QUERY PLAN}

\begin{verbatim}
HashAggregate  (cost=220.00..221.00 rows=100 width=16) (actual rows=100 loops=1)
  Group Key: a, b
    ->  Seq Scan on t  (cost=0.00..145.00 rows=10000 width=8)
         (actual rows=10000 loops=1)
\end{verbatim}

70.3. Planner Statistics and Security

Access to the table \texttt{pg_statistic} is restricted to superusers, so that ordinary users cannot learn about the contents of the tables of other users from it. Some selectivity estimation functions will use a user-provided operator (either the operator appearing in the query or a related operator) to analyze the stored statistics. For example, in order to determine whether a stored most common value is applicable, the selectivity estimator will have to run the appropriate = operator to compare the constant in the query to the stored value. Thus the data in \texttt{pg_statistic} is potentially passed to user-defined operators. An appropriately crafted operator can intentionally leak the passed operands (for example,
by logging them or writing them to a different table), or accidentally leak them by showing their values in error messages, in either case possibly exposing data from \texttt{pg\_statistic} to a user who should not be able to see it.

In order to prevent this, the following applies to all built-in selectivity estimation functions. When planning a query, in order to be able to use stored statistics, the current user must either have \texttt{SELECT} privilege on the table or the involved columns, or the operator used must be \texttt{LEAKPROOF} (more accurately, the function that the operator is based on). If not, then the selectivity estimator will behave as if no statistics are available, and the planner will proceed with default or fall-back assumptions.

If a user does not have the required privilege on the table or columns, then in many cases the query will ultimately receive a permission-denied error, in which case this mechanism is invisible in practice. But if the user is reading from a security-barrier view, then the planner might wish to check the statistics of an underlying table that is otherwise inaccessible to the user. In that case, the operator should be leak-proof or the statistics will not be used. There is no direct feedback about that, except that the plan might be suboptimal. If one suspects that this is the case, one could try running the query as a more privileged user, to see if a different plan results.

This restriction applies only to cases where the planner would need to execute a user-defined operator on one or more values from \texttt{pg\_statistic}. Thus the planner is permitted to use generic statistical information, such as the fraction of null values or the number of distinct values in a column, regardless of access privileges.

Selectivity estimation functions contained in third-party extensions that potentially operate on statistics with user-defined operators should follow the same security rules. Consult the PostgreSQL source code for guidance.
Part VIII. Appendixes
# Table of Contents

A. PostgreSQL Error Codes ................................................................. 2292
B. Date/Time Support .................................................................. 2301
   B.1. Date/Time Input Interpretation ........................................... 2301
   B.2. Date/Time Key Words ......................................................... 2302
   B.3. Date/Time Configuration Files ........................................... 2303
   B.4. History of Units ................................................................. 2304
C. SQL Key Words ......................................................................... 2306
D. SQL Conformance .................................................................... 2329
   D.1. Supported Features ............................................................ 2330
   D.2. Unsupported Features ......................................................... 2345
E. Release Notes ........................................................................... 2358
   E.1. Release 11.1 ................................................................... 2358
      E.1.1. Migration to Version 11.1 ............................................. 2358
      E.1.2. Changes .................................................................. 2358
   E.2. Release 11 ....................................................................... 2360
      E.2.1. Overview .................................................................. 2360
      E.2.2. Migration to Version 11 ............................................. 2361
      E.2.3. Changes .................................................................. 2363
      E.2.4. Acknowledgments ...................................................... 2374
   E.3. Release 10.6 .................................................................... 2379
      E.3.1. Migration to Version 10.6 ............................................. 2379
      E.3.2. Changes .................................................................. 2379
   E.4. Release 10.5 .................................................................... 2384
      E.4.1. Migration to Version 10.5 ............................................. 2385
      E.4.2. Changes .................................................................. 2385
   E.5. Release 10.4 .................................................................... 2389
      E.5.1. Migration to Version 10.4 ............................................. 2389
      E.5.2. Changes .................................................................. 2389
   E.6. Release 10.3 .................................................................... 2393
      E.6.1. Migration to Version 10.3 ............................................. 2393
      E.6.2. Changes .................................................................. 2393
   E.7. Release 10.2 .................................................................... 2395
      E.7.1. Migration to Version 10.2 ............................................. 2395
      E.7.2. Changes .................................................................. 2395
   E.8. Release 10.1 .................................................................... 2400
      E.8.1. Migration to Version 10.1 ............................................. 2400
      E.8.2. Changes .................................................................. 2400
   E.9. Release 10 .................................................................... 2403
      E.9.1. Overview .................................................................. 2403
      E.9.2. Migration to Version 10 .............................................. 2403
      E.9.3. Changes .................................................................. 2406
      E.9.4. Acknowledgments ...................................................... 2418
   E.10. Release 9.6.11 ................................................................. 2423
      E.10.1. Migration to Version 9.6.11 ....................................... 2423
      E.10.2. Changes .................................................................. 2423
   E.11. Release 9.6.10 ................................................................. 2428
      E.11.1. Migration to Version 9.6.10 ....................................... 2428
      E.11.2. Changes .................................................................. 2428
   E.12. Release 9.6.9 ................................................................. 2431
      E.12.1. Migration to Version 9.6.9 ....................................... 2431
      E.12.2. Changes .................................................................. 2431
   E.13. Release 9.6.8 ................................................................. 2434
      E.13.1. Migration to Version 9.6.8 ....................................... 2434
      E.13.2. Changes .................................................................. 2435
   E.14. Release 9.6.7 ................................................................. 2436

2268
E.14.1. Migration to Version 9.6.7 ............................................ 2436
E.14.2. Changes ................................................................. 2436
E.15. Release 9.6.6 .............................................................. 2439
  E.15.1. Migration to Version 9.6.6 ........................................ 2439
  E.15.2. Changes .............................................................. 2439
E.16. Release 9.5.5 .............................................................. 2442
  E.16.1. Migration to Version 9.5.5 ........................................ 2442
  E.16.2. Changes .............................................................. 2442
E.17. Release 9.6.4 .............................................................. 2444
  E.17.1. Migration to Version 9.6.4 ........................................ 2444
  E.17.2. Changes .............................................................. 2444
E.18. Release 9.6.3 .............................................................. 2449
  E.18.1. Migration to Version 9.6.3 ........................................ 2449
  E.18.2. Changes .............................................................. 2449
E.19. Release 9.6.2 .............................................................. 2454
  E.19.1. Migration to Version 9.6.2 ........................................ 2454
  E.19.2. Changes .............................................................. 2454
E.20. Release 9.6.1 .............................................................. 2459
  E.20.1. Migration to Version 9.6.1 ........................................ 2459
  E.20.2. Changes .............................................................. 2459
E.21. Release 9.6 ............................................................... 2462
  E.21.1. Overview ............................................................ 2462
  E.21.2. Migration to Version 9.6 .......................................... 2462
  E.21.3. Changes .............................................................. 2463
E.22. Release 9.5.15 ......................................................... 2480
  E.22.1. Migration to Version 9.5.15 ...................................... 2481
  E.22.2. Changes .............................................................. 2481
E.23. Release 9.5.14 ......................................................... 2484
  E.23.1. Migration to Version 9.5.14 ...................................... 2485
  E.23.2. Changes .............................................................. 2485
E.24. Release 9.5.13 ......................................................... 2487
  E.24.1. Migration to Version 9.5.13 ...................................... 2487
  E.24.2. Changes .............................................................. 2488
E.25. Release 9.5.12 ......................................................... 2490
  E.25.1. Migration to Version 9.5.12 ...................................... 2490
  E.25.2. Changes .............................................................. 2491
E.26. Release 9.5.11 ......................................................... 2492
  E.26.1. Migration to Version 9.5.11 ...................................... 2492
  E.26.2. Changes .............................................................. 2492
E.27. Release 9.5.10 ......................................................... 2495
  E.27.1. Migration to Version 9.5.10 ...................................... 2495
  E.27.2. Changes .............................................................. 2495
E.28. Release 9.5.9 ........................................................... 2497
  E.28.1. Migration to Version 9.5.9 ...................................... 2497
  E.28.2. Changes .............................................................. 2497
E.29. Release 9.5.8 ........................................................... 2499
  E.29.1. Migration to Version 9.5.8 ...................................... 2499
  E.29.2. Changes .............................................................. 2499
E.30. Release 9.5.7 ........................................................... 2504
  E.30.1. Migration to Version 9.5.7 ...................................... 2504
  E.30.2. Changes .............................................................. 2504
E.31. Release 9.5.6 ........................................................... 2508
  E.31.1. Migration to Version 9.5.6 ...................................... 2508
  E.31.2. Changes .............................................................. 2508
E.32. Release 9.5.5 ........................................................... 2512
  E.32.1. Migration to Version 9.5.5 ...................................... 2512
  E.32.2. Changes .............................................................. 2512
E.33. Release 9.5.4 ........................................................... 2516
<table>
<thead>
<tr>
<th>Appendix</th>
<th>Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>E.33.1</td>
<td>Migration to Version 9.5.4</td>
<td>2516</td>
</tr>
<tr>
<td>E.33.2</td>
<td>Changes</td>
<td>2516</td>
</tr>
<tr>
<td>E.34.1</td>
<td>Migration to Version 9.5.3</td>
<td>2520</td>
</tr>
<tr>
<td>E.34.2</td>
<td>Changes</td>
<td>2520</td>
</tr>
<tr>
<td>E.35.1</td>
<td>Migration to Version 9.5.2</td>
<td>2522</td>
</tr>
<tr>
<td>E.35.2</td>
<td>Changes</td>
<td>2522</td>
</tr>
<tr>
<td>E.36.1</td>
<td>Migration to Version 9.5.1</td>
<td>2525</td>
</tr>
<tr>
<td>E.36.2</td>
<td>Changes</td>
<td>2525</td>
</tr>
<tr>
<td>E.37.1</td>
<td>Overview</td>
<td>2527</td>
</tr>
<tr>
<td>E.37.2</td>
<td>Migration to Version 9.5</td>
<td>2527</td>
</tr>
<tr>
<td>E.37.3</td>
<td>Changes</td>
<td>2528</td>
</tr>
<tr>
<td>E.38.1</td>
<td>Migration to Version 9.4.20</td>
<td>2540</td>
</tr>
<tr>
<td>E.38.2</td>
<td>Changes</td>
<td>2540</td>
</tr>
<tr>
<td>E.39.1</td>
<td>Migration to Version 9.4.19</td>
<td>2543</td>
</tr>
<tr>
<td>E.39.2</td>
<td>Changes</td>
<td>2543</td>
</tr>
<tr>
<td>E.40.1</td>
<td>Migration to Version 9.4.18</td>
<td>2546</td>
</tr>
<tr>
<td>E.40.2</td>
<td>Changes</td>
<td>2546</td>
</tr>
<tr>
<td>E.41.1</td>
<td>Migration to Version 9.4.17</td>
<td>2548</td>
</tr>
<tr>
<td>E.41.2</td>
<td>Changes</td>
<td>2548</td>
</tr>
<tr>
<td>E.42.1</td>
<td>Migration to Version 9.4.16</td>
<td>2550</td>
</tr>
<tr>
<td>E.42.2</td>
<td>Changes</td>
<td>2550</td>
</tr>
<tr>
<td>E.43.1</td>
<td>Migration to Version 9.4.15</td>
<td>2552</td>
</tr>
<tr>
<td>E.43.2</td>
<td>Changes</td>
<td>2552</td>
</tr>
<tr>
<td>E.44.1</td>
<td>Migration to Version 9.4.14</td>
<td>2554</td>
</tr>
<tr>
<td>E.44.2</td>
<td>Changes</td>
<td>2554</td>
</tr>
<tr>
<td>E.45.1</td>
<td>Migration to Version 9.4.13</td>
<td>2556</td>
</tr>
<tr>
<td>E.45.2</td>
<td>Changes</td>
<td>2556</td>
</tr>
<tr>
<td>E.46.1</td>
<td>Migration to Version 9.4.12</td>
<td>2561</td>
</tr>
<tr>
<td>E.46.2</td>
<td>Changes</td>
<td>2561</td>
</tr>
<tr>
<td>E.47.1</td>
<td>Migration to Version 9.4.11</td>
<td>2564</td>
</tr>
<tr>
<td>E.47.2</td>
<td>Changes</td>
<td>2565</td>
</tr>
<tr>
<td>E.48.1</td>
<td>Migration to Version 9.4.10</td>
<td>2568</td>
</tr>
<tr>
<td>E.48.2</td>
<td>Changes</td>
<td>2568</td>
</tr>
<tr>
<td>E.49.1</td>
<td>Migration to Version 9.4.9</td>
<td>2571</td>
</tr>
<tr>
<td>E.49.2</td>
<td>Changes</td>
<td>2571</td>
</tr>
<tr>
<td>E.50.1</td>
<td>Migration to Version 9.4.8</td>
<td>2575</td>
</tr>
<tr>
<td>E.50.2</td>
<td>Changes</td>
<td>2575</td>
</tr>
<tr>
<td>E.51.1</td>
<td>Migration to Version 9.4.7</td>
<td>2576</td>
</tr>
<tr>
<td>E.51.2</td>
<td>Changes</td>
<td>2577</td>
</tr>
<tr>
<td>E.52.1</td>
<td>Migration to Version 9.4.6</td>
<td>2578</td>
</tr>
<tr>
<td>Appendix</td>
<td>Description</td>
<td>Page</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>------</td>
</tr>
<tr>
<td>E.52.1.</td>
<td>Migration to Version 9.4.6</td>
<td>2578</td>
</tr>
<tr>
<td>E.52.2.</td>
<td>Changes</td>
<td>2578</td>
</tr>
<tr>
<td>E.53.</td>
<td>Release 9.4.5</td>
<td>2582</td>
</tr>
<tr>
<td>E.53.1.</td>
<td>Migration to Version 9.4.5</td>
<td>2583</td>
</tr>
<tr>
<td>E.53.2.</td>
<td>Changes</td>
<td>2583</td>
</tr>
<tr>
<td>E.54.</td>
<td>Release 9.4.4</td>
<td>2587</td>
</tr>
<tr>
<td>E.54.1.</td>
<td>Migration to Version 9.4.4</td>
<td>2587</td>
</tr>
<tr>
<td>E.54.2.</td>
<td>Changes</td>
<td>2587</td>
</tr>
<tr>
<td>E.55.</td>
<td>Release 9.4.3</td>
<td>2588</td>
</tr>
<tr>
<td>E.55.1.</td>
<td>Migration to Version 9.4.3</td>
<td>2588</td>
</tr>
<tr>
<td>E.55.2.</td>
<td>Changes</td>
<td>2588</td>
</tr>
<tr>
<td>E.56.</td>
<td>Release 9.4.2</td>
<td>2589</td>
</tr>
<tr>
<td>E.56.1.</td>
<td>Migration to Version 9.4.2</td>
<td>2589</td>
</tr>
<tr>
<td>E.56.2.</td>
<td>Changes</td>
<td>2589</td>
</tr>
<tr>
<td>E.57.</td>
<td>Release 9.4.1</td>
<td>2594</td>
</tr>
<tr>
<td>E.57.1.</td>
<td>Migration to Version 9.4.1</td>
<td>2594</td>
</tr>
<tr>
<td>E.57.2.</td>
<td>Changes</td>
<td>2594</td>
</tr>
<tr>
<td>E.58.</td>
<td>Release 9.4</td>
<td>2597</td>
</tr>
<tr>
<td>E.58.1.</td>
<td>Overview</td>
<td>2597</td>
</tr>
<tr>
<td>E.58.2.</td>
<td>Migration to Version 9.4</td>
<td>2597</td>
</tr>
<tr>
<td>E.58.3.</td>
<td>Changes</td>
<td>2600</td>
</tr>
<tr>
<td>E.59.</td>
<td>Release 9.3.25</td>
<td>2611</td>
</tr>
<tr>
<td>E.59.1.</td>
<td>Migration to Version 9.3.25</td>
<td>2612</td>
</tr>
<tr>
<td>E.59.2.</td>
<td>Changes</td>
<td>2612</td>
</tr>
<tr>
<td>E.60.</td>
<td>Release 9.3.24</td>
<td>2614</td>
</tr>
<tr>
<td>E.60.1.</td>
<td>Migration to Version 9.3.24</td>
<td>2615</td>
</tr>
<tr>
<td>E.60.2.</td>
<td>Changes</td>
<td>2615</td>
</tr>
<tr>
<td>E.61.</td>
<td>Release 9.3.23</td>
<td>2617</td>
</tr>
<tr>
<td>E.61.1.</td>
<td>Migration to Version 9.3.23</td>
<td>2617</td>
</tr>
<tr>
<td>E.61.2.</td>
<td>Changes</td>
<td>2617</td>
</tr>
<tr>
<td>E.62.</td>
<td>Release 9.3.22</td>
<td>2619</td>
</tr>
<tr>
<td>E.62.1.</td>
<td>Migration to Version 9.3.22</td>
<td>2619</td>
</tr>
<tr>
<td>E.62.2.</td>
<td>Changes</td>
<td>2619</td>
</tr>
<tr>
<td>E.63.</td>
<td>Release 9.3.21</td>
<td>2620</td>
</tr>
<tr>
<td>E.63.1.</td>
<td>Migration to Version 9.3.21</td>
<td>2621</td>
</tr>
<tr>
<td>E.63.2.</td>
<td>Changes</td>
<td>2621</td>
</tr>
<tr>
<td>E.64.</td>
<td>Release 9.3.20</td>
<td>2623</td>
</tr>
<tr>
<td>E.64.1.</td>
<td>Migration to Version 9.3.20</td>
<td>2623</td>
</tr>
<tr>
<td>E.64.2.</td>
<td>Changes</td>
<td>2623</td>
</tr>
<tr>
<td>E.65.</td>
<td>Release 9.3.19</td>
<td>2624</td>
</tr>
<tr>
<td>E.65.1.</td>
<td>Migration to Version 9.3.19</td>
<td>2624</td>
</tr>
<tr>
<td>E.65.2.</td>
<td>Changes</td>
<td>2624</td>
</tr>
<tr>
<td>E.66.</td>
<td>Release 9.3.18</td>
<td>2626</td>
</tr>
<tr>
<td>E.66.1.</td>
<td>Migration to Version 9.3.18</td>
<td>2626</td>
</tr>
<tr>
<td>E.66.2.</td>
<td>Changes</td>
<td>2626</td>
</tr>
<tr>
<td>E.67.</td>
<td>Release 9.3.17</td>
<td>2630</td>
</tr>
<tr>
<td>E.67.1.</td>
<td>Migration to Version 9.3.17</td>
<td>2630</td>
</tr>
<tr>
<td>E.67.2.</td>
<td>Changes</td>
<td>2630</td>
</tr>
<tr>
<td>E.68.</td>
<td>Release 9.3.16</td>
<td>2633</td>
</tr>
<tr>
<td>E.68.1.</td>
<td>Migration to Version 9.3.16</td>
<td>2633</td>
</tr>
<tr>
<td>E.68.2.</td>
<td>Changes</td>
<td>2633</td>
</tr>
<tr>
<td>E.69.</td>
<td>Release 9.3.15</td>
<td>2636</td>
</tr>
<tr>
<td>E.69.1.</td>
<td>Migration to Version 9.3.15</td>
<td>2636</td>
</tr>
<tr>
<td>E.69.2.</td>
<td>Changes</td>
<td>2636</td>
</tr>
<tr>
<td>E.70.</td>
<td>Release 9.3.14</td>
<td>2638</td>
</tr>
<tr>
<td>E.70.1.</td>
<td>Migration to Version 9.3.14</td>
<td>2639</td>
</tr>
<tr>
<td>E.70.2.</td>
<td>Changes</td>
<td>2639</td>
</tr>
<tr>
<td>E.71.</td>
<td>Release 9.3.13</td>
<td>2642</td>
</tr>
<tr>
<td>Appendix</td>
<td>Title</td>
<td>Page</td>
</tr>
<tr>
<td>----------</td>
<td>-------</td>
<td>------</td>
</tr>
<tr>
<td>E.71.1</td>
<td>Migration to Version 9.3.13</td>
<td>2642</td>
</tr>
<tr>
<td>E.71.2</td>
<td>Changes</td>
<td>2642</td>
</tr>
<tr>
<td>E.72.</td>
<td>Release 9.3.12</td>
<td>2643</td>
</tr>
<tr>
<td>E.72.1</td>
<td>Migration to Version 9.3.12</td>
<td>2643</td>
</tr>
<tr>
<td>E.72.2</td>
<td>Changes</td>
<td>2644</td>
</tr>
<tr>
<td>E.73.</td>
<td>Release 9.3.11</td>
<td>2645</td>
</tr>
<tr>
<td>E.73.1</td>
<td>Migration to Version 9.3.11</td>
<td>2645</td>
</tr>
<tr>
<td>E.73.2</td>
<td>Changes</td>
<td>2645</td>
</tr>
<tr>
<td>E.74.</td>
<td>Release 9.3.10</td>
<td>2649</td>
</tr>
<tr>
<td>E.74.1</td>
<td>Migration to Version 9.3.10</td>
<td>2649</td>
</tr>
<tr>
<td>E.74.2</td>
<td>Changes</td>
<td>2649</td>
</tr>
<tr>
<td>E.75.</td>
<td>Release 9.3.9</td>
<td>2653</td>
</tr>
<tr>
<td>E.75.1</td>
<td>Migration to Version 9.3.9</td>
<td>2653</td>
</tr>
<tr>
<td>E.75.2</td>
<td>Changes</td>
<td>2653</td>
</tr>
<tr>
<td>E.76.</td>
<td>Release 9.3.8</td>
<td>2654</td>
</tr>
<tr>
<td>E.76.1</td>
<td>Migration to Version 9.3.8</td>
<td>2654</td>
</tr>
<tr>
<td>E.76.2</td>
<td>Changes</td>
<td>2654</td>
</tr>
<tr>
<td>E.77.</td>
<td>Release 9.3.7</td>
<td>2655</td>
</tr>
<tr>
<td>E.77.1</td>
<td>Migration to Version 9.3.7</td>
<td>2655</td>
</tr>
<tr>
<td>E.77.2</td>
<td>Changes</td>
<td>2655</td>
</tr>
<tr>
<td>E.78.</td>
<td>Release 9.3.6</td>
<td>2659</td>
</tr>
<tr>
<td>E.78.1</td>
<td>Migration to Version 9.3.6</td>
<td>2659</td>
</tr>
<tr>
<td>E.78.2</td>
<td>Changes</td>
<td>2659</td>
</tr>
<tr>
<td>E.79.</td>
<td>Release 9.3.5</td>
<td>2666</td>
</tr>
<tr>
<td>E.79.1</td>
<td>Migration to Version 9.3.5</td>
<td>2666</td>
</tr>
<tr>
<td>E.79.2</td>
<td>Changes</td>
<td>2667</td>
</tr>
<tr>
<td>E.80.</td>
<td>Release 9.3.4</td>
<td>2671</td>
</tr>
<tr>
<td>E.80.1</td>
<td>Migration to Version 9.3.4</td>
<td>2671</td>
</tr>
<tr>
<td>E.80.2</td>
<td>Changes</td>
<td>2671</td>
</tr>
<tr>
<td>E.81.</td>
<td>Release 9.3.3</td>
<td>2673</td>
</tr>
<tr>
<td>E.81.1</td>
<td>Migration to Version 9.3.3</td>
<td>2673</td>
</tr>
<tr>
<td>E.81.2</td>
<td>Changes</td>
<td>2673</td>
</tr>
<tr>
<td>E.82.</td>
<td>Release 9.3.2</td>
<td>2679</td>
</tr>
<tr>
<td>E.82.1</td>
<td>Migration to Version 9.3.2</td>
<td>2679</td>
</tr>
<tr>
<td>E.82.2</td>
<td>Changes</td>
<td>2679</td>
</tr>
<tr>
<td>E.83.</td>
<td>Release 9.3.1</td>
<td>2682</td>
</tr>
<tr>
<td>E.83.1</td>
<td>Migration to Version 9.3.1</td>
<td>2682</td>
</tr>
<tr>
<td>E.83.2</td>
<td>Changes</td>
<td>2682</td>
</tr>
<tr>
<td>E.84.</td>
<td>Release 9.3</td>
<td>2683</td>
</tr>
<tr>
<td>E.84.1</td>
<td>Overview</td>
<td>2683</td>
</tr>
<tr>
<td>E.84.2</td>
<td>Migration to Version 9.3</td>
<td>2683</td>
</tr>
<tr>
<td>E.84.3</td>
<td>Changes</td>
<td>2684</td>
</tr>
<tr>
<td>E.85.</td>
<td>Release 9.2.24</td>
<td>2694</td>
</tr>
<tr>
<td>E.85.1</td>
<td>Migration to Version 9.2.24</td>
<td>2695</td>
</tr>
<tr>
<td>E.85.2</td>
<td>Changes</td>
<td>2695</td>
</tr>
<tr>
<td>E.86.</td>
<td>Release 9.2.23</td>
<td>2696</td>
</tr>
<tr>
<td>E.86.1</td>
<td>Migration to Version 9.2.23</td>
<td>2696</td>
</tr>
<tr>
<td>E.86.2</td>
<td>Changes</td>
<td>2696</td>
</tr>
<tr>
<td>E.87.</td>
<td>Release 9.2.22</td>
<td>2697</td>
</tr>
<tr>
<td>E.87.1</td>
<td>Migration to Version 9.2.22</td>
<td>2698</td>
</tr>
<tr>
<td>E.87.2</td>
<td>Changes</td>
<td>2698</td>
</tr>
<tr>
<td>E.88.</td>
<td>Release 9.2.21</td>
<td>2701</td>
</tr>
<tr>
<td>E.88.1</td>
<td>Migration to Version 9.2.21</td>
<td>2701</td>
</tr>
<tr>
<td>E.88.2</td>
<td>Changes</td>
<td>2701</td>
</tr>
<tr>
<td>E.89.</td>
<td>Release 9.2.20</td>
<td>2704</td>
</tr>
<tr>
<td>E.89.1</td>
<td>Migration to Version 9.2.20</td>
<td>2704</td>
</tr>
<tr>
<td>E.89.2</td>
<td>Changes</td>
<td>2704</td>
</tr>
<tr>
<td>E.90.</td>
<td>Release 9.2.19</td>
<td>2707</td>
</tr>
</tbody>
</table>
Appendixes

E.90.1. Migration to Version 9.2.19 ......................................................... 2707
E.90.2. Changes ................................................................. 2707
E.91. Release 9.2.18 ................................................................. 2708
   E.91.1. Migration to Version 9.2.18 ....................................... 2709
   E.91.2. Changes ......................................................... 2709
E.92. Release 9.2.17 ................................................................. 2711
   E.92.1. Migration to Version 9.2.17 ................................... 2711
   E.92.2. Changes ......................................................... 2711
E.93. Release 9.2.16 ................................................................. 2712
   E.93.1. Migration to Version 9.2.16 ................................... 2712
   E.93.2. Changes ......................................................... 2712
E.94. Release 9.2.15 ................................................................. 2713
   E.94.1. Migration to Version 9.2.15 ................................... 2714
   E.94.2. Changes ......................................................... 2714
E.95. Release 9.2.14 ................................................................. 2717
   E.95.1. Migration to Version 9.2.14 ................................... 2717
   E.95.2. Changes ......................................................... 2717
E.96. Release 9.2.13 ................................................................. 2721
   E.96.1. Migration to Version 9.2.13 ................................... 2721
   E.96.2. Changes ......................................................... 2721
E.97. Release 9.2.12 ................................................................. 2721
   E.97.1. Migration to Version 9.2.12 ................................... 2721
   E.97.2. Changes ......................................................... 2721
E.98. Release 9.2.11 ................................................................. 2722
   E.98.1. Migration to Version 9.2.11 ................................... 2722
   E.98.2. Changes ......................................................... 2722
E.99. Release 9.2.10 ................................................................. 2726
   E.99.1. Migration to Version 9.2.10 ................................... 2726
   E.99.2. Changes ......................................................... 2726
E.100. Release 9.2.9 ................................................................. 2732
   E.100.1. Migration to Version 9.2.9 ................................... 2732
   E.100.2. Changes ......................................................... 2733
E.101. Release 9.2.8 ................................................................. 2735
   E.101.1. Migration to Version 9.2.8 ................................... 2736
   E.101.2. Changes ......................................................... 2736
E.102. Release 9.2.7 ................................................................. 2737
   E.102.1. Migration to Version 9.2.7 ................................... 2737
   E.102.2. Changes ......................................................... 2737
E.103. Release 9.2.6 ................................................................. 2741
   E.103.1. Migration to Version 9.2.6 ................................... 2741
   E.103.2. Changes ......................................................... 2741
E.104. Release 9.2.5 ................................................................. 2743
   E.104.1. Migration to Version 9.2.5 ................................... 2743
   E.104.2. Changes ......................................................... 2743
E.105. Release 9.2.4 ................................................................. 2746
   E.105.1. Migration to Version 9.2.4 ................................... 2746
   E.105.2. Changes ......................................................... 2746
E.106. Release 9.2.3 ................................................................. 2748
   E.106.1. Migration to Version 9.2.3 ................................... 2748
   E.106.2. Changes ......................................................... 2748
E.107. Release 9.2.2 ................................................................. 2751
   E.107.1. Migration to Version 9.2.2 ................................... 2751
   E.107.2. Changes ......................................................... 2751
E.108. Release 9.2.1 ................................................................. 2756
   E.108.1. Migration to Version 9.2.1 ................................... 2756
   E.108.2. Changes ......................................................... 2756
E.109. Release 9.2 ................................................................. 2757
   E.109.1. Overview ........................................................ 2757

2273
Appendixes

E.109.2. Migration to Version 9.2 .............................................................. 2758
E.109.3. Changes .................................................................................. 2760
E.110. Release 9.1.24 ........................................................................... 2774
E.110.1. Migration to Version 9.1.24 ....................................................... 2774
E.110.2. Changes .................................................................................. 2774
E.111. Release 9.1.23 ........................................................................... 2775
E.111.1. Migration to Version 9.1.23 ....................................................... 2775
E.111.2. Changes .................................................................................. 2775
E.112. Release 9.1.22 ........................................................................... 2778
E.112.1. Migration to Version 9.1.22 ....................................................... 2778
E.112.2. Changes .................................................................................. 2778
E.113. Release 9.1.21 ........................................................................... 2779
E.113.1. Migration to Version 9.1.21 ....................................................... 2779
E.113.2. Changes .................................................................................. 2779
E.114. Release 9.1.20 ........................................................................... 2780
E.114.1. Migration to Version 9.1.20 ....................................................... 2780
E.114.2. Changes .................................................................................. 2780
E.115. Release 9.1.19 ........................................................................... 2783
E.115.1. Migration to Version 9.1.19 ....................................................... 2783
E.115.2. Changes .................................................................................. 2783
E.116. Release 9.1.18 ........................................................................... 2787
E.116.1. Migration to Version 9.1.18 ....................................................... 2787
E.116.2. Changes .................................................................................. 2787
E.117. Release 9.1.17 ........................................................................... 2787
E.117.1. Migration to Version 9.1.17 ....................................................... 2787
E.117.2. Changes .................................................................................. 2788
E.118. Release 9.1.16 ........................................................................... 2788
E.118.1. Migration to Version 9.1.16 ....................................................... 2788
E.118.2. Changes .................................................................................. 2788
E.119. Release 9.1.15 ........................................................................... 2792
E.119.1. Migration to Version 9.1.15 ....................................................... 2792
E.119.2. Changes .................................................................................. 2792
E.120. Release 9.1.14 ........................................................................... 2797
E.120.1. Migration to Version 9.1.14 ....................................................... 2797
E.120.2. Changes .................................................................................. 2798
E.121. Release 9.1.13 ........................................................................... 2800
E.121.1. Migration to Version 9.1.13 ....................................................... 2800
E.121.2. Changes .................................................................................. 2800
E.122. Release 9.1.12 ........................................................................... 2801
E.122.1. Migration to Version 9.1.12 ....................................................... 2801
E.122.2. Changes .................................................................................. 2801
E.123. Release 9.1.11 ........................................................................... 2805
E.123.1. Migration to Version 9.1.11 ....................................................... 2805
E.123.2. Changes .................................................................................. 2805
E.124. Release 9.1.10 ........................................................................... 2806
E.124.1. Migration to Version 9.1.10 ....................................................... 2807
E.124.2. Changes .................................................................................. 2807
E.125. Release 9.1.9 .............................................................................. 2808
E.125.1. Migration to Version 9.1.9 ....................................................... 2809
E.125.2. Changes .................................................................................. 2809
E.126. Release 9.1.8 .............................................................................. 2811
E.126.1. Migration to Version 9.1.8 ....................................................... 2811
E.126.2. Changes .................................................................................. 2811
E.127. Release 9.1.7 .............................................................................. 2813
E.127.1. Migration to Version 9.1.7 ....................................................... 2813
E.127.2. Changes .................................................................................. 2813
E.128. Release 9.1.6 .............................................................................. 2816
E.128.1. Migration to Version 9.1.6 ....................................................... 2816

2274
<table>
<thead>
<tr>
<th>Appendixes</th>
</tr>
</thead>
<tbody>
<tr>
<td>E.147.2. Changes ................................................................. 2876</td>
</tr>
<tr>
<td>E.148. Release 9.0.10 ............................................................. 2878</td>
</tr>
<tr>
<td>E.148.1. Migration to Version 9.0.10 ......................................... 2878</td>
</tr>
<tr>
<td>E.148.2. Changes ..................................................................... 2878</td>
</tr>
<tr>
<td>E.149. Release 9.0.9 ................................................................. 2879</td>
</tr>
<tr>
<td>E.149.1. Migration to Version 9.0.9 ........................................... 2879</td>
</tr>
<tr>
<td>E.149.2. Changes ..................................................................... 2879</td>
</tr>
<tr>
<td>E.150. Release 9.0.8 ................................................................. 2881</td>
</tr>
<tr>
<td>E.150.1. Migration to Version 9.0.8 ........................................... 2881</td>
</tr>
<tr>
<td>E.150.2. Changes ..................................................................... 2881</td>
</tr>
<tr>
<td>E.151. Release 9.0.7 ................................................................. 2883</td>
</tr>
<tr>
<td>E.151.1. Migration to Version 9.0.7 ........................................... 2883</td>
</tr>
<tr>
<td>E.151.2. Changes ..................................................................... 2883</td>
</tr>
<tr>
<td>E.152. Release 9.0.6 ................................................................. 2887</td>
</tr>
<tr>
<td>E.152.1. Migration to Version 9.0.6 ........................................... 2887</td>
</tr>
<tr>
<td>E.152.2. Changes ..................................................................... 2887</td>
</tr>
<tr>
<td>E.153. Release 9.0.5 ................................................................. 2889</td>
</tr>
<tr>
<td>E.153.1. Migration to Version 9.0.5 ........................................... 2889</td>
</tr>
<tr>
<td>E.153.2. Changes ..................................................................... 2889</td>
</tr>
<tr>
<td>E.154. Release 9.0.4 ................................................................. 2893</td>
</tr>
<tr>
<td>E.154.1. Migration to Version 9.0.4 ........................................... 2893</td>
</tr>
<tr>
<td>E.154.2. Changes ..................................................................... 2893</td>
</tr>
<tr>
<td>E.155. Release 9.0.3 ................................................................. 2895</td>
</tr>
<tr>
<td>E.155.1. Migration to Version 9.0.3 ........................................... 2895</td>
</tr>
<tr>
<td>E.155.2. Changes ..................................................................... 2895</td>
</tr>
<tr>
<td>E.156. Release 9.0.2 ................................................................. 2896</td>
</tr>
<tr>
<td>E.156.1. Migration to Version 9.0.2 ........................................... 2897</td>
</tr>
<tr>
<td>E.156.2. Changes ..................................................................... 2897</td>
</tr>
<tr>
<td>E.157. Release 9.0.1 ................................................................. 2899</td>
</tr>
<tr>
<td>E.157.1. Migration to Version 9.0.1 ........................................... 2899</td>
</tr>
<tr>
<td>E.157.2. Changes ..................................................................... 2899</td>
</tr>
<tr>
<td>E.158. Release 9.0 ................................................................. 2900</td>
</tr>
<tr>
<td>E.158.1. Overview ................................................................... 2900</td>
</tr>
<tr>
<td>E.158.2. Migration to Version 9.0 ............................................. 2901</td>
</tr>
<tr>
<td>E.158.3. Changes ..................................................................... 2903</td>
</tr>
<tr>
<td>E.159. Release 8.4.22 ............................................................... 2920</td>
</tr>
<tr>
<td>E.159.1. Migration to Version 8.4.22 ....................................... 2920</td>
</tr>
<tr>
<td>E.159.2. Changes ..................................................................... 2920</td>
</tr>
<tr>
<td>E.160. Release 8.4.21 ............................................................... 2922</td>
</tr>
<tr>
<td>E.160.1. Migration to Version 8.4.21 ....................................... 2922</td>
</tr>
<tr>
<td>E.160.2. Changes ..................................................................... 2922</td>
</tr>
<tr>
<td>E.161. Release 8.4.20 ............................................................... 2923</td>
</tr>
<tr>
<td>E.161.1. Migration to Version 8.4.20 ....................................... 2923</td>
</tr>
<tr>
<td>E.161.2. Changes ..................................................................... 2923</td>
</tr>
<tr>
<td>E.162. Release 8.4.19 ............................................................... 2926</td>
</tr>
<tr>
<td>E.162.1. Migration to Version 8.4.19 ....................................... 2926</td>
</tr>
<tr>
<td>E.162.2. Changes ..................................................................... 2926</td>
</tr>
<tr>
<td>E.163. Release 8.4.18 ............................................................... 2927</td>
</tr>
<tr>
<td>E.163.1. Migration to Version 8.4.18 ....................................... 2927</td>
</tr>
<tr>
<td>E.163.2. Changes ..................................................................... 2928</td>
</tr>
<tr>
<td>E.164. Release 8.4.17 ............................................................... 2929</td>
</tr>
<tr>
<td>E.164.1. Migration to Version 8.4.17 ....................................... 2929</td>
</tr>
<tr>
<td>E.164.2. Changes ..................................................................... 2929</td>
</tr>
<tr>
<td>E.165. Release 8.4.16 ............................................................... 2930</td>
</tr>
<tr>
<td>E.165.1. Migration to Version 8.4.16 ....................................... 2930</td>
</tr>
<tr>
<td>E.165.2. Changes ..................................................................... 2930</td>
</tr>
<tr>
<td>E.166. Release 8.4.15 ............................................................... 2931</td>
</tr>
<tr>
<td>E.166.1. Migration to Version 8.4.15 ....................................... 2932</td>
</tr>
<tr>
<td>Appendix Section</td>
</tr>
<tr>
<td>------------------</td>
</tr>
<tr>
<td>E.205.</td>
</tr>
<tr>
<td>E.205.1.</td>
</tr>
<tr>
<td>E.205.2.</td>
</tr>
<tr>
<td>E.205.3.</td>
</tr>
<tr>
<td>E.206.</td>
</tr>
<tr>
<td>E.206.1.</td>
</tr>
<tr>
<td>E.206.2.</td>
</tr>
<tr>
<td>E.207.</td>
</tr>
<tr>
<td>E.207.1.</td>
</tr>
<tr>
<td>E.207.2.</td>
</tr>
<tr>
<td>E.208.</td>
</tr>
<tr>
<td>E.208.1.</td>
</tr>
<tr>
<td>E.208.2.</td>
</tr>
<tr>
<td>E.209.</td>
</tr>
<tr>
<td>E.209.1.</td>
</tr>
<tr>
<td>E.209.2.</td>
</tr>
<tr>
<td>E.210.</td>
</tr>
<tr>
<td>E.210.1.</td>
</tr>
<tr>
<td>E.210.2.</td>
</tr>
<tr>
<td>E.211.</td>
</tr>
<tr>
<td>E.211.1.</td>
</tr>
<tr>
<td>E.211.2.</td>
</tr>
<tr>
<td>E.212.</td>
</tr>
<tr>
<td>E.212.1.</td>
</tr>
<tr>
<td>E.212.2.</td>
</tr>
<tr>
<td>E.213.</td>
</tr>
<tr>
<td>E.213.1.</td>
</tr>
<tr>
<td>E.213.2.</td>
</tr>
<tr>
<td>E.214.</td>
</tr>
<tr>
<td>E.214.1.</td>
</tr>
<tr>
<td>E.214.2.</td>
</tr>
<tr>
<td>E.215.</td>
</tr>
<tr>
<td>E.215.1.</td>
</tr>
<tr>
<td>E.215.2.</td>
</tr>
<tr>
<td>E.216.</td>
</tr>
<tr>
<td>E.216.1.</td>
</tr>
<tr>
<td>E.216.2.</td>
</tr>
<tr>
<td>E.217.</td>
</tr>
<tr>
<td>E.217.1.</td>
</tr>
<tr>
<td>E.217.2.</td>
</tr>
<tr>
<td>E.218.</td>
</tr>
<tr>
<td>E.218.1.</td>
</tr>
<tr>
<td>E.218.2.</td>
</tr>
<tr>
<td>E.219.</td>
</tr>
<tr>
<td>E.219.1.</td>
</tr>
<tr>
<td>E.219.2.</td>
</tr>
<tr>
<td>E.220.</td>
</tr>
<tr>
<td>E.220.1.</td>
</tr>
<tr>
<td>E.220.2.</td>
</tr>
<tr>
<td>E.221.</td>
</tr>
<tr>
<td>E.221.1.</td>
</tr>
<tr>
<td>E.221.2.</td>
</tr>
<tr>
<td>E.222.</td>
</tr>
<tr>
<td>E.222.1.</td>
</tr>
<tr>
<td>E.222.2.</td>
</tr>
<tr>
<td>E.223.</td>
</tr>
<tr>
<td>E.223.1.</td>
</tr>
<tr>
<td>E.223.2.</td>
</tr>
</tbody>
</table>
Appendixes

E.224. Release 8.2.5 ................................................................. 3063
  E.224.1. Migration to Version 8.2.5 ......................................... 3063
  E.224.2. Changes .............................................................. 3064
E.225. Release 8.2.4 ............................................................. 3064
  E.225.1. Migration to Version 8.2.4 ......................................... 3065
  E.225.2. Changes .............................................................. 3065
E.226. Release 8.2.3 ............................................................. 3065
  E.226.1. Migration to Version 8.2.3 ......................................... 3065
  E.226.2. Changes .............................................................. 3065
E.227. Release 8.2.2 ............................................................. 3066
  E.227.1. Migration to Version 8.2.2 ......................................... 3066
  E.227.2. Changes .............................................................. 3066
E.228. Release 8.2.1 ............................................................. 3067
  E.228.1. Migration to Version 8.2.1 ......................................... 3067
  E.228.2. Changes .............................................................. 3067
E.229. Release 8.2 ............................................................... 3067
  E.229.1. Overview ............................................................. 3067
  E.229.2. Migration to Version 8.2 ........................................... 3068
  E.229.3. Changes .............................................................. 3070
E.230. Release 8.1.23 ......................................................... 3081
  E.230.1. Migration to Version 8.1.23 ........................................ 3082
  E.230.2. Changes .............................................................. 3082
E.231. Release 8.1.22 ......................................................... 3083
  E.231.1. Migration to Version 8.1.22 ........................................ 3083
  E.231.2. Changes .............................................................. 3083
E.232. Release 8.1.21 ......................................................... 3084
  E.232.1. Migration to Version 8.1.21 ........................................ 3084
  E.232.2. Changes .............................................................. 3085
E.233. Release 8.1.20 ......................................................... 3085
  E.233.1. Migration to Version 8.1.20 ........................................ 3086
  E.233.2. Changes .............................................................. 3086
E.234. Release 8.1.19 ......................................................... 3087
  E.234.1. Migration to Version 8.1.19 ........................................ 3087
  E.234.2. Changes .............................................................. 3087
E.235. Release 8.1.18 ......................................................... 3088
  E.235.1. Migration to Version 8.1.18 ........................................ 3088
  E.235.2. Changes .............................................................. 3088
E.236. Release 8.1.17 ......................................................... 3089
  E.236.1. Migration to Version 8.1.17 ........................................ 3089
  E.236.2. Changes .............................................................. 3089
E.237. Release 8.1.16 ......................................................... 3090
  E.237.1. Migration to Version 8.1.16 ........................................ 3090
  E.237.2. Changes .............................................................. 3090
E.238. Release 8.1.15 ......................................................... 3091
  E.238.1. Migration to Version 8.1.15 ........................................ 3091
  E.238.2. Changes .............................................................. 3091
E.239. Release 8.1.14 ......................................................... 3092
  E.239.1. Migration to Version 8.1.14 ........................................ 3092
  E.239.2. Changes .............................................................. 3092
E.240. Release 8.1.13 ......................................................... 3093
  E.240.1. Migration to Version 8.1.13 ........................................ 3093
  E.240.2. Changes .............................................................. 3093
E.241. Release 8.1.12 ......................................................... 3093
  E.241.1. Migration to Version 8.1.12 ........................................ 3093
  E.241.2. Changes .............................................................. 3093
E.242. Release 8.1.11 ......................................................... 3095
  E.242.1. Migration to Version 8.1.11 ........................................ 3095
  E.242.2. Changes .............................................................. 3095
### Appendixes

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>E.300. Release 7.4.11</td>
<td>3173</td>
</tr>
<tr>
<td>E.300.1. Migration to Version 7.4.11</td>
<td>3173</td>
</tr>
<tr>
<td>E.300.2. Changes</td>
<td>3173</td>
</tr>
<tr>
<td>E.301. Release 7.4.10</td>
<td>3173</td>
</tr>
<tr>
<td>E.301.1. Migration to Version 7.4.10</td>
<td>3173</td>
</tr>
<tr>
<td>E.301.2. Changes</td>
<td>3174</td>
</tr>
<tr>
<td>E.302. Release 7.4.9</td>
<td>3174</td>
</tr>
<tr>
<td>E.302.1. Migration to Version 7.4.9</td>
<td>3174</td>
</tr>
<tr>
<td>E.302.2. Changes</td>
<td>3174</td>
</tr>
<tr>
<td>E.303. Release 7.4.8</td>
<td>3175</td>
</tr>
<tr>
<td>E.303.1. Migration to Version 7.4.8</td>
<td>3175</td>
</tr>
<tr>
<td>E.303.2. Changes</td>
<td>3176</td>
</tr>
<tr>
<td>E.304. Release 7.4.7</td>
<td>3177</td>
</tr>
<tr>
<td>E.304.1. Migration to Version 7.4.7</td>
<td>3177</td>
</tr>
<tr>
<td>E.304.2. Changes</td>
<td>3177</td>
</tr>
<tr>
<td>E.305. Release 7.4.6</td>
<td>3178</td>
</tr>
<tr>
<td>E.305.1. Migration to Version 7.4.6</td>
<td>3178</td>
</tr>
<tr>
<td>E.305.2. Changes</td>
<td>3178</td>
</tr>
<tr>
<td>E.306. Release 7.4.5</td>
<td>3179</td>
</tr>
<tr>
<td>E.306.1. Migration to Version 7.4.5</td>
<td>3179</td>
</tr>
<tr>
<td>E.306.2. Changes</td>
<td>3179</td>
</tr>
<tr>
<td>E.307. Release 7.4.4</td>
<td>3179</td>
</tr>
<tr>
<td>E.307.1. Migration to Version 7.4.4</td>
<td>3179</td>
</tr>
<tr>
<td>E.307.2. Changes</td>
<td>3179</td>
</tr>
<tr>
<td>E.308. Release 7.4.3</td>
<td>3180</td>
</tr>
<tr>
<td>E.308.1. Migration to Version 7.4.3</td>
<td>3180</td>
</tr>
<tr>
<td>E.308.2. Changes</td>
<td>3180</td>
</tr>
<tr>
<td>E.309. Release 7.4.2</td>
<td>3181</td>
</tr>
<tr>
<td>E.309.1. Migration to Version 7.4.2</td>
<td>3181</td>
</tr>
<tr>
<td>E.309.2. Changes</td>
<td>3182</td>
</tr>
<tr>
<td>E.310. Release 7.4.1</td>
<td>3182</td>
</tr>
<tr>
<td>E.310.1. Migration to Version 7.4.1</td>
<td>3183</td>
</tr>
<tr>
<td>E.310.2. Changes</td>
<td>3183</td>
</tr>
<tr>
<td>E.311. Release 7.4</td>
<td>3184</td>
</tr>
<tr>
<td>E.311.1. Overview</td>
<td>3184</td>
</tr>
<tr>
<td>E.311.2. Migration to Version 7.4</td>
<td>3186</td>
</tr>
<tr>
<td>E.311.3. Changes</td>
<td>3187</td>
</tr>
<tr>
<td>E.312. Release 7.3.21</td>
<td>3199</td>
</tr>
<tr>
<td>E.312.1. Migration to Version 7.3.21</td>
<td>3199</td>
</tr>
<tr>
<td>E.312.2. Changes</td>
<td>3199</td>
</tr>
<tr>
<td>E.313. Release 7.3.20</td>
<td>3200</td>
</tr>
<tr>
<td>E.313.1. Migration to Version 7.3.20</td>
<td>3200</td>
</tr>
<tr>
<td>E.313.2. Changes</td>
<td>3200</td>
</tr>
<tr>
<td>E.314. Release 7.3.19</td>
<td>3200</td>
</tr>
<tr>
<td>E.314.1. Migration to Version 7.3.19</td>
<td>3200</td>
</tr>
<tr>
<td>E.314.2. Changes</td>
<td>3200</td>
</tr>
<tr>
<td>E.315. Release 7.3.18</td>
<td>3201</td>
</tr>
<tr>
<td>E.315.1. Migration to Version 7.3.18</td>
<td>3201</td>
</tr>
<tr>
<td>E.315.2. Changes</td>
<td>3201</td>
</tr>
<tr>
<td>E.316. Release 7.3.17</td>
<td>3201</td>
</tr>
<tr>
<td>E.316.1. Migration to Version 7.3.17</td>
<td>3201</td>
</tr>
<tr>
<td>E.316.2. Changes</td>
<td>3201</td>
</tr>
<tr>
<td>E.317. Release 7.3.16</td>
<td>3201</td>
</tr>
<tr>
<td>E.317.1. Migration to Version 7.3.16</td>
<td>3201</td>
</tr>
<tr>
<td>E.317.2. Changes</td>
<td>3202</td>
</tr>
<tr>
<td>E.318. Release 7.3.15</td>
<td>3202</td>
</tr>
<tr>
<td>E.318.1. Migration to Version 7.3.15</td>
<td>3202</td>
</tr>
<tr>
<td>E.318.2. Changes</td>
<td>3202</td>
</tr>
</tbody>
</table>

2284
<table>
<thead>
<tr>
<th>Section</th>
<th>Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>E.319.</td>
<td>Release 7.3.14</td>
<td>3203</td>
</tr>
<tr>
<td>E.319.1</td>
<td>Migration to Version 7.3.14</td>
<td>3203</td>
</tr>
<tr>
<td>E.319.2</td>
<td>Changes</td>
<td>3203</td>
</tr>
<tr>
<td>E.320.</td>
<td>Release 7.3.13</td>
<td>3203</td>
</tr>
<tr>
<td>E.320.1</td>
<td>Migration to Version 7.3.13</td>
<td>3203</td>
</tr>
<tr>
<td>E.320.2</td>
<td>Changes</td>
<td>3203</td>
</tr>
<tr>
<td>E.321.</td>
<td>Release 7.3.12</td>
<td>3204</td>
</tr>
<tr>
<td>E.321.1</td>
<td>Migration to Version 7.3.12</td>
<td>3204</td>
</tr>
<tr>
<td>E.321.2</td>
<td>Changes</td>
<td>3204</td>
</tr>
<tr>
<td>E.322.</td>
<td>Release 7.3.11</td>
<td>3204</td>
</tr>
<tr>
<td>E.322.1</td>
<td>Migration to Version 7.3.11</td>
<td>3204</td>
</tr>
<tr>
<td>E.322.2</td>
<td>Changes</td>
<td>3205</td>
</tr>
<tr>
<td>E.323.</td>
<td>Release 7.3.10</td>
<td>3205</td>
</tr>
<tr>
<td>E.323.1</td>
<td>Migration to Version 7.3.10</td>
<td>3205</td>
</tr>
<tr>
<td>E.323.2</td>
<td>Changes</td>
<td>3206</td>
</tr>
<tr>
<td>E.324.</td>
<td>Release 7.3.9</td>
<td>3206</td>
</tr>
<tr>
<td>E.324.1</td>
<td>Migration to Version 7.3.9</td>
<td>3207</td>
</tr>
<tr>
<td>E.324.2</td>
<td>Changes</td>
<td>3207</td>
</tr>
<tr>
<td>E.325.</td>
<td>Release 7.3.8</td>
<td>3207</td>
</tr>
<tr>
<td>E.325.1</td>
<td>Migration to Version 7.3.8</td>
<td>3207</td>
</tr>
<tr>
<td>E.325.2</td>
<td>Changes</td>
<td>3207</td>
</tr>
<tr>
<td>E.326.</td>
<td>Release 7.3.7</td>
<td>3208</td>
</tr>
<tr>
<td>E.326.1</td>
<td>Migration to Version 7.3.7</td>
<td>3208</td>
</tr>
<tr>
<td>E.326.2</td>
<td>Changes</td>
<td>3208</td>
</tr>
<tr>
<td>E.327.</td>
<td>Release 7.3.6</td>
<td>3208</td>
</tr>
<tr>
<td>E.327.1</td>
<td>Migration to Version 7.3.6</td>
<td>3208</td>
</tr>
<tr>
<td>E.327.2</td>
<td>Changes</td>
<td>3208</td>
</tr>
<tr>
<td>E.328.</td>
<td>Release 7.3.5</td>
<td>3209</td>
</tr>
<tr>
<td>E.328.1</td>
<td>Migration to Version 7.3.5</td>
<td>3209</td>
</tr>
<tr>
<td>E.328.2</td>
<td>Changes</td>
<td>3209</td>
</tr>
<tr>
<td>E.329.</td>
<td>Release 7.3.4</td>
<td>3210</td>
</tr>
<tr>
<td>E.329.1</td>
<td>Migration to Version 7.3.4</td>
<td>3210</td>
</tr>
<tr>
<td>E.329.2</td>
<td>Changes</td>
<td>3210</td>
</tr>
<tr>
<td>E.330.</td>
<td>Release 7.3.3</td>
<td>3210</td>
</tr>
<tr>
<td>E.330.1</td>
<td>Migration to Version 7.3.3</td>
<td>3210</td>
</tr>
<tr>
<td>E.330.2</td>
<td>Changes</td>
<td>3210</td>
</tr>
<tr>
<td>E.331.</td>
<td>Release 7.3.2</td>
<td>3213</td>
</tr>
<tr>
<td>E.331.1</td>
<td>Migration to Version 7.3.2</td>
<td>3213</td>
</tr>
<tr>
<td>E.331.2</td>
<td>Changes</td>
<td>3213</td>
</tr>
<tr>
<td>E.332.</td>
<td>Release 7.3.1</td>
<td>3214</td>
</tr>
<tr>
<td>E.332.1</td>
<td>Migration to Version 7.3.1</td>
<td>3214</td>
</tr>
<tr>
<td>E.332.2</td>
<td>Changes</td>
<td>3214</td>
</tr>
<tr>
<td>E.333.</td>
<td>Release 7.3</td>
<td>3214</td>
</tr>
<tr>
<td>E.333.1</td>
<td>Overview</td>
<td>3215</td>
</tr>
<tr>
<td>E.333.2</td>
<td>Migration to Version 7.3</td>
<td>3215</td>
</tr>
<tr>
<td>E.333.3</td>
<td>Changes</td>
<td>3216</td>
</tr>
<tr>
<td>E.334.</td>
<td>Release 7.2.8</td>
<td>3226</td>
</tr>
<tr>
<td>E.334.1</td>
<td>Migration to Version 7.2.8</td>
<td>3226</td>
</tr>
<tr>
<td>E.334.2</td>
<td>Changes</td>
<td>3226</td>
</tr>
<tr>
<td>E.335.</td>
<td>Release 7.2.7</td>
<td>3227</td>
</tr>
<tr>
<td>E.335.1</td>
<td>Migration to Version 7.2.7</td>
<td>3227</td>
</tr>
<tr>
<td>E.335.2</td>
<td>Changes</td>
<td>3227</td>
</tr>
<tr>
<td>E.336.</td>
<td>Release 7.2.6</td>
<td>3227</td>
</tr>
<tr>
<td>E.336.1</td>
<td>Migration to Version 7.2.6</td>
<td>3227</td>
</tr>
<tr>
<td>E.336.2</td>
<td>Changes</td>
<td>3227</td>
</tr>
<tr>
<td>E.337.</td>
<td>Release 7.2.5</td>
<td>3228</td>
</tr>
<tr>
<td>E.337.1</td>
<td>Migration to Version 7.2.5</td>
<td>3228</td>
</tr>
<tr>
<td>E.337.2</td>
<td>Changes</td>
<td>3228</td>
</tr>
</tbody>
</table>
Appendixes

E.338. Release 7.2.4 ................................................................. 3228
E.338.1. Migration to Version 7.2.4 ........................................ 3228
E.338.2. Changes ............................................................... 3228
E.339. Release 7.2.3 ............................................................ 3229
E.339.1. Migration to Version 7.2.3 ....................................... 3229
E.339.2. Changes ............................................................... 3229
E.340. Release 7.2.2 ............................................................ 3229
E.340.1. Migration to Version 7.2.2 ....................................... 3229
E.340.2. Changes ............................................................... 3229
E.341. Release 7.2.1 ............................................................ 3230
E.341.1. Migration to Version 7.2.1 ....................................... 3230
E.341.2. Changes ............................................................... 3230
E.342. Release 7.2 .............................................................. 3231
E.342.1. Overview ............................................................. 3231
E.342.2. Migration to Version 7.2 ......................................... 3231
E.342.3. Changes ............................................................... 3232
E.343. Release 7.1.3 ............................................................ 3241
E.343.1. Migration to Version 7.1.3 ....................................... 3241
E.343.2. Changes ............................................................... 3241
E.344. Release 7.1.2 ............................................................ 3241
E.344.1. Migration to Version 7.1.2 ....................................... 3241
E.344.2. Changes ............................................................... 3241
E.345. Release 7.1.1 ............................................................ 3241
E.345.1. Migration to Version 7.1.1 ....................................... 3241
E.345.2. Changes ............................................................... 3242
E.346. Release 7.1 .............................................................. 3242
E.346.1. Migration to Version 7.1 .......................................... 3243
E.346.2. Changes ............................................................... 3243
E.347. Release 7.0.3 ............................................................ 3246
E.347.1. Migration to Version 7.0.3 ....................................... 3246
E.347.2. Changes ............................................................... 3246
E.348. Release 7.0.2 ............................................................ 3247
E.348.1. Migration to Version 7.0.2 ....................................... 3247
E.348.2. Changes ............................................................... 3247
E.349. Release 7.0.1 ............................................................ 3248
E.349.1. Migration to Version 7.0.1 ....................................... 3248
E.349.2. Changes ............................................................... 3248
E.350. Release 7.0 .............................................................. 3248
E.350.1. Migration to Version 7.0 .......................................... 3249
E.350.2. Changes ............................................................... 3249
E.351. Release 6.5.3 ............................................................ 3256
E.351.1. Migration to Version 6.5.3 ....................................... 3256
E.351.2. Changes ............................................................... 3256
E.352. Release 6.5.2 ............................................................ 3256
E.352.1. Migration to Version 6.5.2 ....................................... 3256
E.352.2. Changes ............................................................... 3256
E.353. Release 6.5.1 ............................................................ 3257
E.353.1. Migration to Version 6.5.1 ....................................... 3257
E.353.2. Changes ............................................................... 3257
E.354. Release 6.5 .............................................................. 3257
E.354.1. Migration to Version 6.5 .......................................... 3258
E.354.2. Changes ............................................................... 3259
E.355. Release 6.4.2 ............................................................ 3262
E.355.1. Migration to Version 6.4.2 ....................................... 3262
E.355.2. Changes ............................................................... 3263
E.356. Release 6.4.1 ............................................................ 3263
E.356.1. Migration to Version 6.4.1 ....................................... 3263
E.356.2. Changes ............................................................... 3263
<table>
<thead>
<tr>
<th>Section</th>
<th>Description</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>F.5.2</td>
<td>Examples</td>
<td>3301</td>
</tr>
<tr>
<td>F.5.3</td>
<td>Operator Class Interface</td>
<td>3303</td>
</tr>
<tr>
<td>F.5.4</td>
<td>Limitations</td>
<td>3303</td>
</tr>
<tr>
<td>F.5.5</td>
<td>Authors</td>
<td>3304</td>
</tr>
<tr>
<td>F.6</td>
<td>btree_gin</td>
<td>3304</td>
</tr>
<tr>
<td>F.6.1</td>
<td>Example Usage</td>
<td>3304</td>
</tr>
<tr>
<td>F.6.2</td>
<td>Authors</td>
<td>3304</td>
</tr>
<tr>
<td>F.7</td>
<td>btree_gist</td>
<td>3304</td>
</tr>
<tr>
<td>F.7.1</td>
<td>Example Usage</td>
<td>3305</td>
</tr>
<tr>
<td>F.7.2</td>
<td>Authors</td>
<td>3305</td>
</tr>
<tr>
<td>F.8</td>
<td>citext</td>
<td>3305</td>
</tr>
<tr>
<td>F.8.1</td>
<td>Rationale</td>
<td>3306</td>
</tr>
<tr>
<td>F.8.2</td>
<td>How to Use It</td>
<td>3306</td>
</tr>
<tr>
<td>F.8.3</td>
<td>String Comparison Behavior</td>
<td>3306</td>
</tr>
<tr>
<td>F.8.4</td>
<td>Limitations</td>
<td>3307</td>
</tr>
<tr>
<td>F.8.5</td>
<td>Author</td>
<td>3307</td>
</tr>
<tr>
<td>F.9</td>
<td>cube</td>
<td>3308</td>
</tr>
<tr>
<td>F.9.1</td>
<td>Syntax</td>
<td>3308</td>
</tr>
<tr>
<td>F.9.2</td>
<td>Precision</td>
<td>3308</td>
</tr>
<tr>
<td>F.9.3</td>
<td>Usage</td>
<td>3308</td>
</tr>
<tr>
<td>F.9.4</td>
<td>Defaults</td>
<td>3311</td>
</tr>
<tr>
<td>F.9.5</td>
<td>Notes</td>
<td>3312</td>
</tr>
<tr>
<td>F.9.6</td>
<td>Credits</td>
<td>3312</td>
</tr>
<tr>
<td>F.10</td>
<td>dblink</td>
<td>3313</td>
</tr>
<tr>
<td>F.11</td>
<td>dict_int</td>
<td>3344</td>
</tr>
<tr>
<td>F.11.1</td>
<td>Configuration</td>
<td>3344</td>
</tr>
<tr>
<td>F.11.2</td>
<td>Usage</td>
<td>3344</td>
</tr>
<tr>
<td>F.12</td>
<td>dict_xsyn</td>
<td>3345</td>
</tr>
<tr>
<td>F.12.1</td>
<td>Configuration</td>
<td>3345</td>
</tr>
<tr>
<td>F.12.2</td>
<td>Usage</td>
<td>3345</td>
</tr>
<tr>
<td>F.13</td>
<td>earthdistance</td>
<td>3346</td>
</tr>
<tr>
<td>F.13.1</td>
<td>Cube-based Earth Distances</td>
<td>3346</td>
</tr>
<tr>
<td>F.13.2</td>
<td>Point-based Earth Distances</td>
<td>3347</td>
</tr>
<tr>
<td>F.14</td>
<td>file_fdw</td>
<td>3348</td>
</tr>
<tr>
<td>F.15</td>
<td>fuzzystrmatch</td>
<td>3350</td>
</tr>
<tr>
<td>F.15.1</td>
<td>Soundex</td>
<td>3350</td>
</tr>
<tr>
<td>F.15.2</td>
<td>Levenshtein</td>
<td>3351</td>
</tr>
<tr>
<td>F.15.3</td>
<td>Metaphone</td>
<td>3352</td>
</tr>
<tr>
<td>F.15.4</td>
<td>Double Metaphone</td>
<td>3352</td>
</tr>
<tr>
<td>F.16</td>
<td>hstore</td>
<td>3353</td>
</tr>
<tr>
<td>F.16.1</td>
<td>hstore External Representation</td>
<td>3353</td>
</tr>
<tr>
<td>F.16.2</td>
<td>hstore Operators and Functions</td>
<td>3353</td>
</tr>
<tr>
<td>F.16.3</td>
<td>Indexes</td>
<td>3357</td>
</tr>
<tr>
<td>F.16.4</td>
<td>Examples</td>
<td>3357</td>
</tr>
<tr>
<td>F.16.5</td>
<td>Statistics</td>
<td>3358</td>
</tr>
<tr>
<td>F.16.6</td>
<td>Compatibility</td>
<td>3359</td>
</tr>
<tr>
<td>F.16.7</td>
<td>Transforms</td>
<td>3359</td>
</tr>
<tr>
<td>F.16.8</td>
<td>Authors</td>
<td>3359</td>
</tr>
<tr>
<td>F.17</td>
<td>intagg</td>
<td>3359</td>
</tr>
<tr>
<td>F.17.1</td>
<td>Functions</td>
<td>3359</td>
</tr>
<tr>
<td>F.17.2</td>
<td>Sample Uses</td>
<td>3360</td>
</tr>
<tr>
<td>F.18</td>
<td>intarray</td>
<td>3361</td>
</tr>
<tr>
<td>F.18.1</td>
<td>intarray Functions and Operators</td>
<td>3361</td>
</tr>
<tr>
<td>F.18.2</td>
<td>Index Support</td>
<td>3362</td>
</tr>
<tr>
<td>F.18.3</td>
<td>Example</td>
<td>3362</td>
</tr>
<tr>
<td>F.18.4</td>
<td>Benchmark</td>
<td>3363</td>
</tr>
<tr>
<td>F.18.5</td>
<td>Authors</td>
<td>3363</td>
</tr>
<tr>
<td>F.19</td>
<td>isn</td>
<td>3363</td>
</tr>
</tbody>
</table>
### Appendixes

<table>
<thead>
<tr>
<th>Section</th>
<th>Title</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>F.19.1.</td>
<td>Data Types</td>
<td>3363</td>
</tr>
<tr>
<td>F.19.2.</td>
<td>Casts</td>
<td>3364</td>
</tr>
<tr>
<td>F.19.3.</td>
<td>Functions and Operators</td>
<td>3365</td>
</tr>
<tr>
<td>F.19.4.</td>
<td>Examples</td>
<td>3365</td>
</tr>
<tr>
<td>F.19.5.</td>
<td>Bibliography</td>
<td>3366</td>
</tr>
<tr>
<td>F.19.6.</td>
<td>Author</td>
<td>3367</td>
</tr>
<tr>
<td>F.20.</td>
<td>io</td>
<td>3367</td>
</tr>
<tr>
<td>F.20.1.</td>
<td>Rationale</td>
<td>3367</td>
</tr>
<tr>
<td>F.20.2.</td>
<td>How to Use It</td>
<td>3367</td>
</tr>
<tr>
<td>F.20.3.</td>
<td>Limitations</td>
<td>3367</td>
</tr>
<tr>
<td>F.20.4.</td>
<td>Author</td>
<td>3368</td>
</tr>
<tr>
<td>F.21.</td>
<td>ltree</td>
<td>3368</td>
</tr>
<tr>
<td>F.21.1.</td>
<td>Definitions</td>
<td>3368</td>
</tr>
<tr>
<td>F.21.2.</td>
<td>Operators and Functions</td>
<td>3369</td>
</tr>
<tr>
<td>F.21.3.</td>
<td>Indexes</td>
<td>3372</td>
</tr>
<tr>
<td>F.21.4.</td>
<td>Example</td>
<td>3372</td>
</tr>
<tr>
<td>F.21.5.</td>
<td>Transforms</td>
<td>3374</td>
</tr>
<tr>
<td>F.21.6.</td>
<td>Authors</td>
<td>3374</td>
</tr>
<tr>
<td>F.22.</td>
<td>pageinspect</td>
<td>3374</td>
</tr>
<tr>
<td>F.22.1.</td>
<td>General Functions</td>
<td>3374</td>
</tr>
<tr>
<td>F.22.2.</td>
<td>B-tree Functions</td>
<td>3376</td>
</tr>
<tr>
<td>F.22.3.</td>
<td>BRIN Functions</td>
<td>3378</td>
</tr>
<tr>
<td>F.22.4.</td>
<td>GIN Functions</td>
<td>3379</td>
</tr>
<tr>
<td>F.22.5.</td>
<td>Hash Functions</td>
<td>3380</td>
</tr>
<tr>
<td>F.23.</td>
<td>passwordcheck</td>
<td>3381</td>
</tr>
<tr>
<td>F.24.</td>
<td>pg_buffercache</td>
<td>3382</td>
</tr>
<tr>
<td>F.24.1.</td>
<td>The pg_buffercache View</td>
<td>3382</td>
</tr>
<tr>
<td>F.24.2.</td>
<td>Sample Output</td>
<td>3383</td>
</tr>
<tr>
<td>F.24.3.</td>
<td>Authors</td>
<td>3383</td>
</tr>
<tr>
<td>F.25.</td>
<td>pgcrypto</td>
<td>3384</td>
</tr>
<tr>
<td>F.25.1.</td>
<td>General Hashing Functions</td>
<td>3384</td>
</tr>
<tr>
<td>F.25.2.</td>
<td>Password Hashing Functions</td>
<td>3384</td>
</tr>
<tr>
<td>F.25.3.</td>
<td>PGP Encryption Functions</td>
<td>3386</td>
</tr>
<tr>
<td>F.25.4.</td>
<td>Raw Encryption Functions</td>
<td>3391</td>
</tr>
<tr>
<td>F.25.5.</td>
<td>Random-Data Functions</td>
<td>3392</td>
</tr>
<tr>
<td>F.25.6.</td>
<td>Notes</td>
<td>3393</td>
</tr>
<tr>
<td>F.25.7.</td>
<td>Author</td>
<td>3394</td>
</tr>
<tr>
<td>F.26.</td>
<td>pg_freespacemap</td>
<td>3395</td>
</tr>
<tr>
<td>F.26.1.</td>
<td>Functions</td>
<td>3395</td>
</tr>
<tr>
<td>F.26.2.</td>
<td>Sample Output</td>
<td>3395</td>
</tr>
<tr>
<td>F.26.3.</td>
<td>Author</td>
<td>3396</td>
</tr>
<tr>
<td>F.27.</td>
<td>pg_prewarm</td>
<td>3396</td>
</tr>
<tr>
<td>F.27.1.</td>
<td>Functions</td>
<td>3396</td>
</tr>
<tr>
<td>F.27.2.</td>
<td>Configuration Parameters</td>
<td>3397</td>
</tr>
<tr>
<td>F.27.3.</td>
<td>Author</td>
<td>3397</td>
</tr>
<tr>
<td>F.28.</td>
<td>pgrowlocks</td>
<td>3397</td>
</tr>
<tr>
<td>F.28.1.</td>
<td>Overview</td>
<td>3397</td>
</tr>
<tr>
<td>F.28.2.</td>
<td>Sample Output</td>
<td>3398</td>
</tr>
<tr>
<td>F.28.3.</td>
<td>Author</td>
<td>3398</td>
</tr>
<tr>
<td>F.29.</td>
<td>pg_stat_statements</td>
<td>3398</td>
</tr>
<tr>
<td>F.29.1.</td>
<td>The pg_stat_statements View</td>
<td>3398</td>
</tr>
<tr>
<td>F.29.2.</td>
<td>Functions</td>
<td>3401</td>
</tr>
<tr>
<td>F.29.3.</td>
<td>Configuration Parameters</td>
<td>3401</td>
</tr>
<tr>
<td>F.29.4.</td>
<td>Sample Output</td>
<td>3402</td>
</tr>
<tr>
<td>F.29.5.</td>
<td>Authors</td>
<td>3403</td>
</tr>
<tr>
<td>F.30.</td>
<td>pgstattuple</td>
<td>3403</td>
</tr>
<tr>
<td>F.30.1.</td>
<td>Functions</td>
<td>3403</td>
</tr>
<tr>
<td>F.30.2.</td>
<td>Authors</td>
<td>3407</td>
</tr>
</tbody>
</table>
Appendices

F.43.3. Functions ................................................................. 3449
F.44. uuid-ossip ................................................................. 3449
   F.44.1. uuid-ossip Functions ........................................ 3449
   F.44.2. Building uuid-ossip ............................................ 3450
   F.44.3. Author .............................................................. 3451
F.45. xml2 ................................................................. 3451
   F.45.1. Deprecation Notice ............................................ 3451
   F.45.2. Description of Functions .................................... 3451
   F.45.3. xpath_table .................................................... 3452
   F.45.4. XSLT Functions ............................................... 3454
   F.45.5. Author ............................................................ 3455
G. Additional Supplied Programs ........................................... 3456
   G.1. Client Applications ............................................... 3456
   G.2. Server Applications ............................................... 3462
H. External Projects ........................................................ 3467
   H.1. Client Interfaces .................................................. 3467
   H.2. Administration Tools ............................................. 3467
   H.3. Procedural Languages ............................................ 3468
   H.4. Extensions .......................................................... 3468
I. The Source Code Repository ............................................. 3469
   I.1. Getting The Source via Git ....................................... 3469
J. Documentation ........................................................... 3470
   J.1. DocBook ............................................................. 3470
   J.2. Tool Sets ............................................................. 3470
      J.2.1. Installation on Fedora, RHEL, and Derivatives ........ 3471
      J.2.2. Installation on FreeBSD ..................................... 3471
      J.2.3. Debian Packages ............................................. 3471
      J.2.4. macOS ......................................................... 3471
      J.2.5. Detection by configure ..................................... 3472
   J.3. Building The Documentation .................................... 3472
      J.3.1. HTML ............................................................ 3472
      J.3.2. Manpages ...................................................... 3472
      J.3.3. PDF ............................................................. 3472
      J.3.4. Plain Text Files ............................................. 3473
      J.3.5. Syntax Check .................................................. 3473
   J.4. Documentation Authoring ........................................ 3473
      J.4.1. Emacs .......................................................... 3473
   J.5. Style Guide ........................................................ 3474
      J.5.1. Reference Pages ............................................... 3474
K. Acronyms ..................................................................... 3476
Appendix A. PostgreSQL Error Codes

All messages emitted by the PostgreSQL server are assigned five-character error codes that follow
the SQL standard’s conventions for “SQLSTATE” codes. Applications that need to know which error
condition has occurred should usually test the error code, rather than looking at the textual error mes-
sage. The error codes are less likely to change across PostgreSQL releases, and also are not subject to
change due to localization of error messages. Note that some, but not all, of the error codes produced
by PostgreSQL are defined by the SQL standard; some additional error codes for conditions not de-
defined by the standard have been invented or borrowed from other databases.

According to the standard, the first two characters of an error code denote a class of errors, while the
last three characters indicate a specific condition within that class. Thus, an application that does not
recognize the specific error code might still be able to infer what to do from the error class.

Table A.1 lists all the error codes defined in PostgreSQL 11.1. (Some are not actually used at present,
but are defined by the SQL standard.) The error classes are also shown. For each error class there is a
“standard” error code having the last three characters 000. This code is used only for error conditions
that fall within the class but do not have any more-specific code assigned.

The symbol shown in the column “Condition Name” is the condition name to use in PL/pgSQL. Con-
dition names can be written in either upper or lower case. (Note that PL/pgSQL does not recognize
warning, as opposed to error, condition names; those are classes 00, 01, and 02.)

For some types of errors, the server reports the name of a database object (a table, table column,
data type, or constraint) associated with the error; for example, the name of the unique constraint that
caused a unique_violation error. Such names are supplied in separate fields of the error report
message so that applications need not try to extract them from the possibly-localized human-readable
text of the message. As of PostgreSQL 9.3, complete coverage for this feature exists only for errors
in SQLSTATE class 23 (integrity constraint violation), but this is likely to be expanded in future.

Table A.1. PostgreSQL Error Codes

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Class 00 — Successful Completion</strong></td>
<td></td>
</tr>
<tr>
<td>00000</td>
<td>successful_completion</td>
</tr>
<tr>
<td><strong>Class 01 — Warning</strong></td>
<td></td>
</tr>
<tr>
<td>01000</td>
<td>warning</td>
</tr>
<tr>
<td>0100C</td>
<td>dynamic_result_sets_returned</td>
</tr>
<tr>
<td>01008</td>
<td>implicit_zero_bit_padding</td>
</tr>
<tr>
<td>01003</td>
<td>null_value_eliminated_in_set_function</td>
</tr>
<tr>
<td>01007</td>
<td>privilege_not_granted</td>
</tr>
<tr>
<td>01006</td>
<td>privilege_not_revoked</td>
</tr>
<tr>
<td>01004</td>
<td>string_data_right_truncation</td>
</tr>
<tr>
<td>01P01</td>
<td>deprecated_feature</td>
</tr>
<tr>
<td><strong>Class 02 — No Data (this is also a warning class per the SQL standard)</strong></td>
<td></td>
</tr>
<tr>
<td>02000</td>
<td>no_data</td>
</tr>
<tr>
<td>02001</td>
<td>no_additional_dynamic_result_sets_returned</td>
</tr>
<tr>
<td><strong>Class 03 — SQL Statement Not Yet Complete</strong></td>
<td></td>
</tr>
<tr>
<td>03000</td>
<td>sql_statement_not_yet_complete</td>
</tr>
<tr>
<td><strong>Class 08 — Connection Exception</strong></td>
<td></td>
</tr>
</tbody>
</table>
### PostgreSQL Error Codes

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>08000</td>
<td>connection_exception</td>
</tr>
<tr>
<td>08003</td>
<td>connection_does_not_exist</td>
</tr>
<tr>
<td>08006</td>
<td>connection_failure</td>
</tr>
<tr>
<td>08001</td>
<td>sqlclient_unable_to_estabish_sqlconnection</td>
</tr>
<tr>
<td>08004</td>
<td>sqlserverRejectedEstablishment_of_sqlconnection</td>
</tr>
<tr>
<td>08007</td>
<td>transaction_resolution_unknown</td>
</tr>
<tr>
<td>08P01</td>
<td>protocol_violation</td>
</tr>
</tbody>
</table>

**Class 09 — Triggered Action Exception**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>09000</td>
<td>triggered_action_exception</td>
</tr>
</tbody>
</table>

**Class 0A — Feature Not Supported**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0A000</td>
<td>feature_not_supported</td>
</tr>
</tbody>
</table>

**Class 0B — Invalid Transaction Initiation**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0B000</td>
<td>invalid_transaction_initiation</td>
</tr>
</tbody>
</table>

**Class 0F — Locator Exception**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0F000</td>
<td>locator_exception</td>
</tr>
<tr>
<td>0F001</td>
<td>invalid_locator_specification</td>
</tr>
</tbody>
</table>

**Class 0L — Invalid Grantor**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0L000</td>
<td>invalid_grantor</td>
</tr>
<tr>
<td>0LP01</td>
<td>invalid_grant_operation</td>
</tr>
</tbody>
</table>

**Class 0P — Invalid Role Specification**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0P000</td>
<td>invalid_role_specification</td>
</tr>
</tbody>
</table>

**Class 0Z — Diagnostics Exception**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>0Z000</td>
<td>diagnostics_exception</td>
</tr>
<tr>
<td>0Z002</td>
<td>stacked_diagnostics_accessed_without_active_handler</td>
</tr>
</tbody>
</table>

**Class 20 — Case Not Found**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>20000</td>
<td>case_not_found</td>
</tr>
</tbody>
</table>

**Class 21 — Cardinality Violation**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>21000</td>
<td>cardinality_violation</td>
</tr>
</tbody>
</table>

**Class 22 — Data Exception**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>22000</td>
<td>data_exception</td>
</tr>
<tr>
<td>2202E</td>
<td>array_subscript_error</td>
</tr>
<tr>
<td>22021</td>
<td>character_not_in_repertoire</td>
</tr>
<tr>
<td>22008</td>
<td>datetime_field_overflow</td>
</tr>
<tr>
<td>22012</td>
<td>division_by_zero</td>
</tr>
<tr>
<td>22005</td>
<td>error_in_assignment</td>
</tr>
<tr>
<td>2200B</td>
<td>escape_character_conflict</td>
</tr>
<tr>
<td>22022</td>
<td>indicator_overflow</td>
</tr>
<tr>
<td>22015</td>
<td>interval_field_overflow</td>
</tr>
<tr>
<td>2201E</td>
<td>invalid_argument_for_logarithm</td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>----------------------------------------------------------</td>
</tr>
<tr>
<td>22014</td>
<td>invalid_argument_for_ntile_function</td>
</tr>
<tr>
<td>22016</td>
<td>invalid_argument_for_nth_value_function</td>
</tr>
<tr>
<td>2201F</td>
<td>invalid_argument_for_power_function</td>
</tr>
<tr>
<td>2201G</td>
<td>invalid_argument_for_width_bucket_function</td>
</tr>
<tr>
<td>22018</td>
<td>invalid_character_value_for_cast</td>
</tr>
<tr>
<td>22007</td>
<td>invalid_datetime_format</td>
</tr>
<tr>
<td>22019</td>
<td>invalid_escape_character</td>
</tr>
<tr>
<td>2200D</td>
<td>invalid_escape_octet</td>
</tr>
<tr>
<td>22025</td>
<td>invalid_escape_sequence</td>
</tr>
<tr>
<td>22P06</td>
<td>nonstandard_use_of_escape_character</td>
</tr>
<tr>
<td>22010</td>
<td>invalid_indicator_parameter_value</td>
</tr>
<tr>
<td>22023</td>
<td>invalid_parameter_value</td>
</tr>
<tr>
<td>22013</td>
<td>invalid_preceding_or_following_size</td>
</tr>
<tr>
<td>2201B</td>
<td>invalid_regular_expression</td>
</tr>
<tr>
<td>2201W</td>
<td>invalid_row_count_in_limit_clause</td>
</tr>
<tr>
<td>2201X</td>
<td>invalid_row_count_in_result_offset_clause</td>
</tr>
<tr>
<td>2202H</td>
<td>invalid_tablesample_argument</td>
</tr>
<tr>
<td>2202G</td>
<td>invalid_tablesample_repeat</td>
</tr>
<tr>
<td>22009</td>
<td>invalid_time_zone_displacement_value</td>
</tr>
<tr>
<td>2200C</td>
<td>invalid_use_of_escape_character</td>
</tr>
<tr>
<td>2200G</td>
<td>most_specific_type_mismatch</td>
</tr>
<tr>
<td>22004</td>
<td>null_value_not_allowed</td>
</tr>
<tr>
<td>22002</td>
<td>null_value_no_indicator_parameter</td>
</tr>
<tr>
<td>22003</td>
<td>numeric_value_out_of_range</td>
</tr>
<tr>
<td>2200H</td>
<td>sequence_generator_limit_exceeded</td>
</tr>
<tr>
<td>22026</td>
<td>string_data_length_mismatch</td>
</tr>
<tr>
<td>22001</td>
<td>string_data_right_truncation</td>
</tr>
<tr>
<td>22011</td>
<td>substring_error</td>
</tr>
<tr>
<td>22027</td>
<td>trim_error</td>
</tr>
<tr>
<td>22024</td>
<td>unterminated_c_string</td>
</tr>
<tr>
<td>2200F</td>
<td>zero_length_character_string</td>
</tr>
<tr>
<td>22P01</td>
<td>floating_point_exception</td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>-----------------------------------------------</td>
</tr>
<tr>
<td>22P02</td>
<td>invalid_text_representation</td>
</tr>
<tr>
<td>22P03</td>
<td>invalid_binary_representation</td>
</tr>
<tr>
<td>22P04</td>
<td>bad_copy_file_format</td>
</tr>
<tr>
<td>22P05</td>
<td>untranslatable_character</td>
</tr>
<tr>
<td>2200L</td>
<td>not_an_xml_document</td>
</tr>
<tr>
<td>2200M</td>
<td>invalid_xml_document</td>
</tr>
<tr>
<td>2200N</td>
<td>invalid_xml_content</td>
</tr>
<tr>
<td>2200S</td>
<td>invalid_xml_comment</td>
</tr>
<tr>
<td>2200T</td>
<td>invalid_xml_processing_instruction</td>
</tr>
</tbody>
</table>

**Class 23 — Integrity Constraint Violation**

| 23000      | integrity_constraintViolation                  |
| 23001      | restrict_violation                             |
| 23502      | not_null_violation                             |
| 23503      | foreign_key_violation                          |
| 23505      | unique_violation                               |
| 23514      | check_violation                                |
| 23P01      | exclusion_violation                            |

**Class 24 — Invalid Cursor State**

| 24000      | invalid_cursor_state                           |

**Class 25 — Invalid Transaction State**

| 25000      | invalid_transaction_state                     |
| 25001      | active_sql_transaction                        |
| 25002      | branch_transaction_already_active             |
| 25008      | held_cursor.requires_same_isolation_level     |
| 25003      | inappropriate_access_mode_for_branch_transaction |
| 25004      | inappropriate_isolation_level_for_branch_transaction |
| 25005      | no_active_sql_transaction_for_branch_transaction |
| 25006      | read_only_sql_transaction                      |
| 25007      | schema_and_data_statement_mixing_not_supported |
| 25P01      | no_active_sql_transaction                      |
| 25P02      | in_failed_sql_transaction                      |
| 25P03      | idle_in_transaction_session_timeout            |

**Class 26 — Invalid SQL Statement Name**

| 26000      | invalid_sql_statement_name                     |

**Class 27 — Triggered Data Change Violation**
<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>27000</td>
<td>triggered_data_change_violation</td>
</tr>
<tr>
<td><strong>Class 28 — Invalid Authorization Specification</strong></td>
<td></td>
</tr>
<tr>
<td>28000</td>
<td>invalid_authorization_specification</td>
</tr>
<tr>
<td>28P01</td>
<td>invalid_password</td>
</tr>
<tr>
<td><strong>Class 2B — Dependent Privilege Descriptors Still Exist</strong></td>
<td></td>
</tr>
<tr>
<td>2B000</td>
<td>dependent_privilege_descriptors_still_exist</td>
</tr>
<tr>
<td>2BP01</td>
<td>dependent_objects_still_exist</td>
</tr>
<tr>
<td><strong>Class 2D — Invalid Transaction Termination</strong></td>
<td></td>
</tr>
<tr>
<td>2D000</td>
<td>invalid_transaction_termination</td>
</tr>
<tr>
<td><strong>Class 2F — SQL Routine Exception</strong></td>
<td></td>
</tr>
<tr>
<td>2F000</td>
<td>sql_routine_exception</td>
</tr>
<tr>
<td>2F005</td>
<td>function_executed_no_return_statement</td>
</tr>
<tr>
<td>2F002</td>
<td>modifying_sql_data_not_permitted</td>
</tr>
<tr>
<td>2F003</td>
<td>prohibited_sql_statement_attempted</td>
</tr>
<tr>
<td>2F004</td>
<td>reading_sql_data_not_permitted</td>
</tr>
<tr>
<td><strong>Class 34 — Invalid Cursor Name</strong></td>
<td></td>
</tr>
<tr>
<td>34000</td>
<td>invalid_cursor_name</td>
</tr>
<tr>
<td><strong>Class 38 — External Routine Exception</strong></td>
<td></td>
</tr>
<tr>
<td>38000</td>
<td>external_routine_exception</td>
</tr>
<tr>
<td>38001</td>
<td>containing_sql_not_permitted</td>
</tr>
<tr>
<td>38002</td>
<td>modifying_sql_data_not_permitted</td>
</tr>
<tr>
<td>38003</td>
<td>prohibited_sql_statement_attempted</td>
</tr>
<tr>
<td>38004</td>
<td>reading_sql_data_not_permitted</td>
</tr>
<tr>
<td><strong>Class 39 — External Routine Invocation Exception</strong></td>
<td></td>
</tr>
<tr>
<td>39000</td>
<td>external_routine_invocation_exception</td>
</tr>
<tr>
<td>39001</td>
<td>invalid_sqlstate_returned</td>
</tr>
<tr>
<td>39004</td>
<td>null_value_not_allowed</td>
</tr>
<tr>
<td>39P01</td>
<td>trigger_protocol_violated</td>
</tr>
<tr>
<td>39P02</td>
<td>srf_protocol_violated</td>
</tr>
<tr>
<td>39P03</td>
<td>event_trigger_protocol_violated</td>
</tr>
<tr>
<td><strong>Class 3B — Savepoint Exception</strong></td>
<td></td>
</tr>
<tr>
<td>3B000</td>
<td>savepoint_exception</td>
</tr>
<tr>
<td>3B001</td>
<td>invalid_savepoint_specification</td>
</tr>
<tr>
<td><strong>Class 3D — Invalid Catalog Name</strong></td>
<td></td>
</tr>
<tr>
<td>3D000</td>
<td>invalid_catalog_name</td>
</tr>
<tr>
<td><strong>Class 3F — Invalid Schema Name</strong></td>
<td></td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>----------------------------------------------------</td>
</tr>
<tr>
<td>3F000</td>
<td>invalid_schema_name</td>
</tr>
</tbody>
</table>

### Class 40 — Transaction Rollback

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>40000</td>
<td>transaction_rollback</td>
</tr>
<tr>
<td>40002</td>
<td>transaction_integrity_constraint_violation</td>
</tr>
<tr>
<td>40001</td>
<td>serialization_failure</td>
</tr>
<tr>
<td>40003</td>
<td>statement_completion_unknown</td>
</tr>
<tr>
<td>40P01</td>
<td>deadlock_detected</td>
</tr>
</tbody>
</table>

### Class 42 — Syntax Error or Access Rule Violation

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>42000</td>
<td>syntax_error_or_access_rule_violation</td>
</tr>
<tr>
<td>42601</td>
<td>syntax_error</td>
</tr>
<tr>
<td>42501</td>
<td>insufficient_privilege</td>
</tr>
<tr>
<td>42846</td>
<td>cannot_coerce</td>
</tr>
<tr>
<td>42803</td>
<td>grouping_error</td>
</tr>
<tr>
<td>42P20</td>
<td>windowing_error</td>
</tr>
<tr>
<td>42P19</td>
<td>invalid_recursion</td>
</tr>
<tr>
<td>42830</td>
<td>invalid_foreign_key</td>
</tr>
<tr>
<td>42602</td>
<td>invalid_name</td>
</tr>
<tr>
<td>42622</td>
<td>name_too_long</td>
</tr>
<tr>
<td>42939</td>
<td>reserved_name</td>
</tr>
<tr>
<td>42804</td>
<td>datatype_mismatch</td>
</tr>
<tr>
<td>42P18</td>
<td>indeterminate_datatype</td>
</tr>
<tr>
<td>42P21</td>
<td>collation_mismatch</td>
</tr>
<tr>
<td>42P22</td>
<td>indeterminate_collation</td>
</tr>
<tr>
<td>42809</td>
<td>wrong_object_type</td>
</tr>
<tr>
<td>428C9</td>
<td>generated_always</td>
</tr>
<tr>
<td>42703</td>
<td>undefined_column</td>
</tr>
<tr>
<td>42883</td>
<td>undefined_function</td>
</tr>
<tr>
<td>42P01</td>
<td>undefined_table</td>
</tr>
<tr>
<td>42P02</td>
<td>undefined_parameter</td>
</tr>
<tr>
<td>42704</td>
<td>undefined_object</td>
</tr>
<tr>
<td>42701</td>
<td>duplicate_column</td>
</tr>
<tr>
<td>42P03</td>
<td>duplicate_cursor</td>
</tr>
<tr>
<td>42P04</td>
<td>duplicate_database</td>
</tr>
<tr>
<td>42723</td>
<td>duplicate_function</td>
</tr>
<tr>
<td>42P05</td>
<td>duplicate_prepared_statement</td>
</tr>
<tr>
<td>42P06</td>
<td>duplicate_schema</td>
</tr>
<tr>
<td>42P07</td>
<td>duplicate_table</td>
</tr>
<tr>
<td>42712</td>
<td>duplicate_alias</td>
</tr>
<tr>
<td>42710</td>
<td>duplicate_object</td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>-----------</td>
<td>------------------------------------</td>
</tr>
<tr>
<td>42702</td>
<td>ambiguous_column</td>
</tr>
<tr>
<td>42725</td>
<td>ambiguous_function</td>
</tr>
<tr>
<td>42P08</td>
<td>ambiguous_parameter</td>
</tr>
<tr>
<td>42P09</td>
<td>ambiguous_alias</td>
</tr>
<tr>
<td>42P10</td>
<td>invalid_column_reference</td>
</tr>
<tr>
<td>42611</td>
<td>invalid_column_definition</td>
</tr>
<tr>
<td>42P11</td>
<td>invalid_cursor_definition</td>
</tr>
<tr>
<td>42P12</td>
<td>invalid_database_definition</td>
</tr>
<tr>
<td>42P13</td>
<td>invalid_function_definition</td>
</tr>
<tr>
<td>42P14</td>
<td>invalid_prepared_statement_definition</td>
</tr>
<tr>
<td>42P15</td>
<td>invalid_schema_definition</td>
</tr>
<tr>
<td>42P16</td>
<td>invalid_table_definition</td>
</tr>
<tr>
<td>42P17</td>
<td>invalid_object_definition</td>
</tr>
</tbody>
</table>

**Class 44 — WITH CHECK OPTION Violation**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>44000</td>
<td>with_check_option_violation</td>
</tr>
</tbody>
</table>

**Class 53 — Insufficient Resources**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>53000</td>
<td>insufficient_resources</td>
</tr>
<tr>
<td>53100</td>
<td>disk_full</td>
</tr>
<tr>
<td>53200</td>
<td>out_of_memory</td>
</tr>
<tr>
<td>53300</td>
<td>too_many_connections</td>
</tr>
<tr>
<td>53400</td>
<td>configuration_limit_exceeded</td>
</tr>
</tbody>
</table>

**Class 54 — Program Limit Exceeded**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>54000</td>
<td>program_limit_exceeded</td>
</tr>
<tr>
<td>54001</td>
<td>statement_too_complex</td>
</tr>
<tr>
<td>54011</td>
<td>too_many_columns</td>
</tr>
<tr>
<td>54023</td>
<td>too_many_arguments</td>
</tr>
</tbody>
</table>

**Class 55 — Object Not In Prerequisite State**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>55000</td>
<td>object_not_in_prerequisite_state</td>
</tr>
<tr>
<td>55006</td>
<td>object_in_use</td>
</tr>
<tr>
<td>55P02</td>
<td>cant_change_runtime_param</td>
</tr>
<tr>
<td>55P03</td>
<td>lock_not_available</td>
</tr>
</tbody>
</table>

**Class 57 — Operator Intervention**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>57000</td>
<td>operator_intervention</td>
</tr>
<tr>
<td>57014</td>
<td>query_canceled</td>
</tr>
<tr>
<td>57P01</td>
<td>admin_shutdown</td>
</tr>
<tr>
<td>57P02</td>
<td>crash_shutdown</td>
</tr>
<tr>
<td>57P03</td>
<td>cannot_connect_now</td>
</tr>
<tr>
<td>57P04</td>
<td>database_dropped</td>
</tr>
</tbody>
</table>

**Class 58 — System Error (errors external to PostgreSQL itself)**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>58000</td>
<td>system_error</td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>------------------------------</td>
</tr>
<tr>
<td>58030</td>
<td>io_error</td>
</tr>
<tr>
<td>58P01</td>
<td>undefined_file</td>
</tr>
<tr>
<td>58P02</td>
<td>duplicate_file</td>
</tr>
</tbody>
</table>

**Class 72 — Snapshot Failure**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>72000</td>
<td>snapshot_too_old</td>
</tr>
</tbody>
</table>

**Class F0 — Configuration File Error**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>F0000</td>
<td>config_file_error</td>
</tr>
<tr>
<td>F0001</td>
<td>lock_file_exists</td>
</tr>
</tbody>
</table>

**Class HV — Foreign Data Wrapper Error (SQL/MED)**

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>HV000</td>
<td>fdw_error</td>
</tr>
<tr>
<td>HV005</td>
<td>fdw_column_name_not_found</td>
</tr>
<tr>
<td>HV002</td>
<td>fdw_dynamic_parameter_value_needed</td>
</tr>
<tr>
<td>HV010</td>
<td>fdw_function_sequence_error</td>
</tr>
<tr>
<td>HV021</td>
<td>fdw_inconsistent_descriptor_information</td>
</tr>
<tr>
<td>HV024</td>
<td>fdw_invalid_attribute_value</td>
</tr>
<tr>
<td>HV007</td>
<td>fdw_invalid_column_name</td>
</tr>
<tr>
<td>HV008</td>
<td>fdw_invalid_column_number</td>
</tr>
<tr>
<td>HV004</td>
<td>fdw_invalid_data_type</td>
</tr>
<tr>
<td>HV006</td>
<td>fdw_invalid_data_type_descriptors</td>
</tr>
<tr>
<td>HV091</td>
<td>fdw_invalid_descriptor_field_identifier</td>
</tr>
<tr>
<td>HV00B</td>
<td>fdw_invalid_handle</td>
</tr>
<tr>
<td>HV00C</td>
<td>fdw_invalid_option_index</td>
</tr>
<tr>
<td>HV00D</td>
<td>fdw_invalid_option_name</td>
</tr>
<tr>
<td>HV090</td>
<td>fdw_invalid_string_length_or_buffer_length</td>
</tr>
<tr>
<td>HV00A</td>
<td>fdw_invalid_string_format</td>
</tr>
<tr>
<td>HV009</td>
<td>fdw_invalid_use_of_null_pointer</td>
</tr>
<tr>
<td>HV014</td>
<td>fdw_too_many_handles</td>
</tr>
<tr>
<td>HV001</td>
<td>fdw_out_of_memory</td>
</tr>
<tr>
<td>HV00P</td>
<td>fdw_no_schemas</td>
</tr>
<tr>
<td>HV00J</td>
<td>fdw_option_name_not_found</td>
</tr>
<tr>
<td>HV00K</td>
<td>fdw_reply_handle</td>
</tr>
<tr>
<td>HV00Q</td>
<td>fdw_schema_not_found</td>
</tr>
<tr>
<td>HV00R</td>
<td>fdw_table_not_found</td>
</tr>
<tr>
<td>HV00L</td>
<td>fdw_unable_to_create_execution</td>
</tr>
<tr>
<td>HV00M</td>
<td>fdw_unable_to_create_reply</td>
</tr>
<tr>
<td>HV00N</td>
<td>fdw_unable_to_establish_connection</td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>------------------------</td>
</tr>
<tr>
<td><strong>Class P0 — PL/pgSQL Error</strong></td>
<td></td>
</tr>
<tr>
<td>P0000</td>
<td>plpgsql_error</td>
</tr>
<tr>
<td>P0001</td>
<td>raise_exception</td>
</tr>
<tr>
<td>P0002</td>
<td>no_data_found</td>
</tr>
<tr>
<td>P0003</td>
<td>too_many_rows</td>
</tr>
<tr>
<td>P0004</td>
<td>assert_failure</td>
</tr>
<tr>
<td><strong>Class XX — Internal Error</strong></td>
<td></td>
</tr>
<tr>
<td>XX000</td>
<td>internal_error</td>
</tr>
<tr>
<td>XX001</td>
<td>data_corrupted</td>
</tr>
<tr>
<td>XX002</td>
<td>index_corrupted</td>
</tr>
</tbody>
</table>
Appendix B. Date/Time Support

PostgreSQL uses an internal heuristic parser for all date/time input support. Dates and times are input as strings, and are broken up into distinct fields with a preliminary determination of what kind of information can be in the field. Each field is interpreted and either assigned a numeric value, ignored, or rejected. The parser contains internal lookup tables for all textual fields, including months, days of the week, and time zones.

This appendix includes information on the content of these lookup tables and describes the steps used by the parser to decode dates and times.

B.1. Date/Time Input Interpretation

The date/time type inputs are all decoded using the following procedure.

1. Break the input string into tokens and categorize each token as a string, time, time zone, or number.
   a. If the numeric token contains a colon (:) this is a time string. Include all subsequent digits and colons.
   b. If the numeric token contains a dash (-), slash (/), or two or more dots (.), this is a date string which might have a text month. If a date token has already been seen, it is instead interpreted as a time zone name (e.g., America/New_York).
   c. If the token is numeric only, then it is either a single field or an ISO 8601 concatenated date (e.g., 19990113 for January 13, 1999) or time (e.g., 141516 for 14:15:16).
   d. If the token starts with a plus (+) or minus (-), then it is either a numeric time zone or a special field.

2. If the token is a text string, match up with possible strings:
   a. Do a binary-search table lookup for the token as a time zone abbreviation.
   b. If not found, do a similar binary-search table lookup to match the token as either a special string (e.g., today), day (e.g., Thursday), month (e.g., January), or noise word (e.g., at, on).
   c. If still not found, throw an error.

3. When the token is a number or number field:
   a. If there are eight or six digits, and if no other date fields have been previously read, then interpret as a “concatenated date” (e.g., 19990118 or 990118). The interpretation is YYYY-MM-DD or YYMMDD.
   b. If the token is three digits and a year has already been read, then interpret as day of year.
   c. If four or six digits and a year has already been read, then interpret as a time (HHMM or HHMMSS).
   d. If three or more digits and no date fields have yet been found, interpret as a year (this forces yy-mm-dd ordering of the remaining date fields).
   e. Otherwise the date field ordering is assumed to follow the DateStyle setting: mm-dd-yyyy, dd-mm-yyyy, or yy-mm-dd. Throw an error if a month or day field is found to be out of range.

4. If BC has been specified, negate the year and add one for internal storage. (There is no year zero in the Gregorian calendar, so numerically 1 BC becomes year zero.)
5. If BC was not specified, and if the year field was two digits in length, then adjust the year to four
digits. If the field is less than 70, then add 2000, otherwise add 1900.

**Tip**

Gregorian years AD 1-99 can be entered by using 4 digits with leading zeros (e.g.,
0099 is AD 99).

### B.2. Date/Time Key Words

Table B.1 shows the tokens that are recognized as names of months.

**Table B.1. Month Names**

<table>
<thead>
<tr>
<th>Month</th>
<th>Abbreviations</th>
</tr>
</thead>
<tbody>
<tr>
<td>January</td>
<td>Jan</td>
</tr>
<tr>
<td>February</td>
<td>Feb</td>
</tr>
<tr>
<td>March</td>
<td>Mar</td>
</tr>
<tr>
<td>April</td>
<td>Apr</td>
</tr>
<tr>
<td>May</td>
<td></td>
</tr>
<tr>
<td>June</td>
<td>Jun</td>
</tr>
<tr>
<td>July</td>
<td>Jul</td>
</tr>
<tr>
<td>August</td>
<td>Aug</td>
</tr>
<tr>
<td>September</td>
<td>Sep, Sept</td>
</tr>
<tr>
<td>October</td>
<td>Oct</td>
</tr>
<tr>
<td>November</td>
<td>Nov</td>
</tr>
<tr>
<td>December</td>
<td>Dec</td>
</tr>
</tbody>
</table>

Table B.2 shows the tokens that are recognized as names of days of the week.

**Table B.2. Day of the Week Names**

<table>
<thead>
<tr>
<th>Day</th>
<th>Abbreviations</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sunday</td>
<td>Sun</td>
</tr>
<tr>
<td>Monday</td>
<td>Mon</td>
</tr>
<tr>
<td>Tuesday</td>
<td>Tue, Tues</td>
</tr>
<tr>
<td>Wednesday</td>
<td>Wed, Weds</td>
</tr>
<tr>
<td>Thursday</td>
<td>Thu, Thur, Thurs</td>
</tr>
<tr>
<td>Friday</td>
<td>Fri</td>
</tr>
<tr>
<td>Saturday</td>
<td>Sat</td>
</tr>
</tbody>
</table>

Table B.3 shows the tokens that serve various modifier purposes.

**Table B.3. Date/Time Field Modifiers**

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AM</td>
<td>Time is before 12:00</td>
</tr>
</tbody>
</table>
Date/Time Support

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AT</td>
<td>Ignored</td>
</tr>
<tr>
<td>JULIAN, JD, J</td>
<td>Next field is Julian Date</td>
</tr>
<tr>
<td>ON</td>
<td>Ignored</td>
</tr>
<tr>
<td>PM</td>
<td>Time is on or after 12:00</td>
</tr>
<tr>
<td>T</td>
<td>Next field is time</td>
</tr>
</tbody>
</table>

### B.3. Date/Time Configuration Files

Since timezone abbreviations are not well standardized, PostgreSQL provides a means to customize the set of abbreviations accepted by the server. The timezone_abbreviations run-time parameter determines the active set of abbreviations. While this parameter can be altered by any database user, the possible values for it are under the control of the database administrator — they are in fact names of configuration files stored in `.../share/timezonesets/` of the installation directory. By adding or altering files in that directory, the administrator can set local policy for timezone abbreviations.

`timezone_abbreviations` can be set to any file name found in `.../share/timezonesets/`, if the file's name is entirely alphabetic. (The prohibition against non-alphabetic characters in `timezone_abbreviations` prevents reading files outside the intended directory, as well as reading editor backup files and other extraneous files.)

A timezone abbreviation file can contain blank lines and comments beginning with `#`. Non-comment lines must have one of these formats:

- `zone_abbreviation offset`
- `zone_abbreviation offset D`
- `zone_abbreviation time_zone_name` @INCLUDE file_name @OVERRIDE

A `zone_abbreviation` is just the abbreviation being defined. An `offset` is an integer giving the equivalent offset in seconds from UTC, positive being east from Greenwich and negative being west. For example, -18000 would be five hours west of Greenwich, or North American east coast standard time. `D` indicates that the zone name represents local daylight-savings time rather than standard time.

Alternatively, a `time_zone_name` can be given, referencing a zone name defined in the IANA timezone database. The zone's definition is consulted to see whether the abbreviation is or has been in use in that zone, and if so, the appropriate meaning is used — that is, the meaning that was currently in use at the timestamp whose value is being determined, or the meaning in use immediately before that if it wasn't current at that time, or the oldest meaning if it was used only after that time. This behavior is essential for dealing with abbreviations whose meaning has historically varied. It is also allowed to define an abbreviation in terms of a zone name in which that abbreviation does not appear; then using the abbreviation is just equivalent to writing out the zone name.

**Tip**

Using a simple integer `offset` is preferred when defining an abbreviation whose offset from UTC has never changed, as such abbreviations are much cheaper to process than those that require consulting a time zone definition.

The `@INCLUDE` syntax allows inclusion of another file in the `.../share/timezonesets/` directory. Inclusion can be nested, to a limited depth.
The \@OVERRIDE syntax indicates that subsequent entries in the file can override previous entries (typically, entries obtained from included files). Without this, conflicting definitions of the same time-zone abbreviation are considered an error.

In an unmodified installation, the file Default contains all the non-conflicting time zone abbreviations for most of the world. Additional files Australia and India are provided for those regions: these files first include the Default file and then add or modify abbreviations as needed.

For reference purposes, a standard installation also contains files Africa.txt, America.txt, etc., containing information about every time zone abbreviation known to be in use according to the IANA timezone database. The zone name definitions found in these files can be copied and pasted into a custom configuration file as needed. Note that these files cannot be directly referenced as time-zone_abbreviations settings, because of the dot embedded in their names.

**Note**

If an error occurs while reading the time zone abbreviation set, no new value is applied and the old set is kept. If the error occurs while starting the database, startup fails.

**Caution**

Time zone abbreviations defined in the configuration file override non-timezone meanings built into PostgreSQL. For example, the Australia configuration file defines SAT (for South Australian Standard Time). When this file is active, SAT will not be recognized as an abbreviation for Saturday.

**Caution**

If you modify files in .../share/timezonesets/, it is up to you to make backups — a normal database dump will not include this directory.

**B.4. History of Units**

The SQL standard states that “Within the definition of a ‘datetime literal’, the ‘datetime values’ are constrained by the natural rules for dates and times according to the Gregorian calendar”. PostgreSQL follows the SQL standard’s lead by counting dates exclusively in the Gregorian calendar, even for years before that calendar was in use. This rule is known as the proleptic Gregorian calendar.

The Julian calendar was introduced by Julius Caesar in 45 BC. It was in common use in the Western world until the year 1582, when countries started changing to the Gregorian calendar. In the Julian calendar, the tropical year is approximated as 365 1/4 days = 365.25 days. This gives an error of about 1 day in 128 years.

The accumulating calendar error prompted Pope Gregory XIII to reform the calendar in accordance with instructions from the Council of Trent. In the Gregorian calendar, the tropical year is approximated as 365 + 97 / 400 days = 365.2425 days. Thus it takes approximately 3300 years for the tropical year to shift one day with respect to the Gregorian calendar.

The approximation 365+97/400 is achieved by having 97 leap years every 400 years, using the following rules:

Every year divisible by 4 is a leap year.
However, every year divisible by 100 is not a leap year.
However, every year divisible by 400 is a leap year after all.

So, 1700, 1800, 1900, 2100, and 2200 are not leap years. But 1600, 2000, and 2400 are leap years. By
contrast, in the older Julian calendar all years divisible by 4 are leap years.

The papal bull of February 1582 decreed that 10 days should be dropped from October 1582 so that
15 October should follow immediately after 4 October. This was observed in Italy, Poland, Portugal,
and Spain. Other Catholic countries followed shortly after, but Protestant countries were reluctant
to change, and the Greek Orthodox countries didn’t change until the start of the 20th century. The
reform was observed by Great Britain and its dominions (including what is now the USA) in 1752.
Thus 2 September 1752 was followed by 14 September 1752. This is why Unix systems have the cal
program produce the following:

```
$ cal 9 1752
  September 1752
  Su Mo Tu We Th Fr Sa
  1  2 14 15 16
  17 18 19 20 21 22 23
  24 25 26 27 28 29 30
```

But, of course, this calendar is only valid for Great Britain and dominions, not other places. Since it
would be difficult and confusing to try to track the actual calendars that were in use in various places
at various times, PostgreSQL does not try, but rather follows the Gregorian calendar rules for all dates,
even though this method is not historically accurate.

Different calendars have been developed in various parts of the world, many predating the Gregorian
system. For example, the beginnings of the Chinese calendar can be traced back to the 14th century
BC. Legend has it that the Emperor Huangdi invented that calendar in 2637 BC. The People’s Republic
of China uses the Gregorian calendar for civil purposes. The Chinese calendar is used for determining
festivals.

The Julian Date system is another type of calendar, unrelated to the Julian calendar though it is con-
fusingly named similarly to that calendar. The Julian Date system was invented by the French schol-
ar Joseph Justus Scaliger (1540-1609) and probably takes its name from Scaliger’s father, the Italian
scholar Julius Caesar Scaliger (1484-1558). In the Julian Date system, each day has a sequential num-
ber, starting from JD 0 (which is sometimes called the Julian Date). JD 0 corresponds to 1 January
4713 BC in the Julian calendar, or 24 November 4714 BC in the Gregorian calendar. Julian Date
counting is most often used by astronomers for labeling their nightly observations, and therefore a date
runs from noon UTC to the next noon UTC, rather than from midnight to midnight: JD 0 designates
the 24 hours from noon UTC on 24 November 4714 BC to noon UTC on 25 November 4714 BC.

Although PostgreSQL supports Julian Date notation for input and output of dates (and also uses Julian
dates for some internal datetime calculations), it does not observe the nicety of having dates run from
noon to noon. PostgreSQL treats a Julian Date as running from midnight to midnight.
Appendix C. SQL Key Words

Table C.1 lists all tokens that are key words in the SQL standard and in PostgreSQL 11.1. Background information can be found in Section 4.1.1. (For space reasons, only the latest two versions of the SQL standard, and SQL-92 for historical comparison, are included. The differences between those and the other intermediate standard versions are small.)

SQL distinguishes between reserved and non-reserved key words. According to the standard, reserved key words are the only real key words; they are never allowed as identifiers. Non-reserved key words only have a special meaning in particular contexts and can be used as identifiers in other contexts. Most non-reserved key words are actually the names of built-in tables and functions specified by SQL. The concept of non-reserved key words essentially only exists to declare that some predefined meaning is attached to a word in some contexts.

In the PostgreSQL parser life is a bit more complicated. There are several different classes of tokens ranging from those that can never be used as an identifier to those that have absolutely no special status in the parser as compared to an ordinary identifier. (The latter is usually the case for functions specified by SQL.) Even reserved key words are not completely reserved in PostgreSQL, but can be used as column labels (for example, SELECT 55 AS CHECK, even though CHECK is a reserved key word).

In Table C.1 in the column for PostgreSQL we classify as “non-reserved” those key words that are explicitly known to the parser but are allowed as column or table names. Some key words that are otherwise non-reserved cannot be used as function or data type names and are marked accordingly. (Most of these words represent built-in functions or data types with special syntax. The function or type is still available but it cannot be redefined by the user.) Labeled “reserved” are those tokens that are not allowed as column or table names. Some reserved key words are allowable as names for functions or data types; this is also shown in the table. If not so marked, a reserved key word is only allowed as an “AS” column label name.

As a general rule, if you get spurious parser errors for commands that contain any of the listed key words as an identifier you should try to quote the identifier to see if the problem goes away.

It is important to understand before studying Table C.1 that the fact that a key word is not reserved in PostgreSQL does not mean that the feature related to the word is not implemented. Conversely, the presence of a key word does not indicate the existence of a feature.

Table C.1. SQL Key Words

<table>
<thead>
<tr>
<th>Key Word</th>
<th>PostgreSQL</th>
<th>SQL:2011</th>
<th>SQL:2008</th>
<th>SQL-92</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABORT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABS</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABSENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ABSOLUTE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ACCESS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ACCORDING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ACTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ADA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ADD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ADMIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>AFTER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>AGGREGATE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>ALLOCATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALSO</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALTER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALWAYS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ANALYSE</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ANALYZE</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>AND</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ANY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARRAY</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ARRAY_AGG</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ARRAY_MAX_CARDINALITY</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>AS</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASC</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASENSITIVE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ASSERTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASSIGNMENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASYMMETRIC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>AT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ATOMIC</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ATTACH</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ATTRIBUTE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>ATTRIBUTES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>AUTHORIZATION</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>AVG</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>BACKWARD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>BASE64</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BEFORE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BEGIN</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>BEGIN_FRAME</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>BEGIN_PARTITION</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>BERNOUlli</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BETWEEN</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>BIGINT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>BINARY</td>
<td>reserved (can be reserved function or type)</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>BIT</td>
<td>non-reserved (cannot be function or type)</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>BIT_LENGTH</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>BLOB</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>BLOCKED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BOM</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BOOLEAN</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>BOTH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>BREADTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>BY</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>C</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CACHE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CALL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CALLED</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CARDINALITY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CASCADE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CASCADED</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CASE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CAST</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CATALOG_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CEIL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CEILING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CHAIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>CHAR</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CHARACTER</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CHARACTERISTICS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>CHARACTERS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>CHARACTER_LENGTH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CHARACTER_SET_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>CHARACTER_SET_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>CHARACTER_SET_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>CHAR_LENGTH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CHECK</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CHECKPOINT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CLASS</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CLASS_ORIGIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CLOB</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CLOSE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CLUSTER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COALESCE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COBOL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLLATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COLLATION</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COLLOCATION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLLOCATION_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLLOCATION_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLLOCATION_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLLECT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COLUMN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COLUMNS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COLUMN_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COMMAND_FUNCTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COMMAND_FUNCTION_CODE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>COMMENT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>COMMENTS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>COMMIT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COMMITTED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONCURRENTLY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONDITION</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONDITION_NUMBER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONFIGURATION</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CONFLICT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CONNECT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL:92</td>
</tr>
<tr>
<td>----------------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>----------</td>
</tr>
<tr>
<td>CONNECTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONNECTION_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONSTRAINT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONSTRAINTS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONSTRAINT_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONSTRAINT_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONSTRAINT_T_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONSTRUCTOR</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONTENT</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONTENTS</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONTINUE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONTROL</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CONVERSION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CONVERT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COPY</td>
<td>reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CORR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CORRESPONDING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COST</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COUNT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COVAR_POP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>COVAR_SAMP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CREATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CROSS</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CSV</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CUBE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CUME_DIST</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_CATALOG</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_DATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_DEFAULT_TRANSFORM_GROUP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>CURRENT_PATH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_ROLE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_ROW</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_T_SCHEMA</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>-----------------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>CURRENT_TIME</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_TIMESTAMP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURRENT_TRANSFORM_GROUP_FOR_TYPE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>CURRENT_USER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURSOR</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>CURSOR_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>CYCLE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DATA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>DATABASE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DATALINK</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DATETIME_INTERVAL_CODE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>DATETIME_INTERVAL_PRECISION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>DAY</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DB</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DEALLOCATE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DEC</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DECIMAL</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DECLARE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DEFAULT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DEFAULTS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>DEFERRABLE</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DEFERRED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DEFINED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>DEFINER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>DEGREE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DELETE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DELIMITER</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DELIMITERS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DENSE_RANK</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DEPENDS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DEPTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DEREF</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>DERIVED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>DESC</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DESCRIBE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DESCRIPTOR</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DETACH</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DETERMINISTIC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DIAGNOSTICS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DICTIONARY</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DISABLE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DISCARD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DISCONNECT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DISPATCH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>DISTINCT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLNEWCOPY</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLPREVIOUSCOPY</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLCOMPLETE</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLCOMPLETEONLY</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLCOMPLETEREWRITE</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLPATH</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLPATHONLY</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLPATHWRITE</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLSCHEME</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLURLSERVER</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DLVALUE</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DO</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DOCUMENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DOMAIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DOUBLE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DROP</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DYNAMIC</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DYNAMIC_FUNCTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>DYNAMIC_FUNCTION_CODE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>EACH</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ELEMENT</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>ELSE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EMPTY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ENABLE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ENCODING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ENCRYPTED</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>END</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>END-EXEC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>END_FRAME</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>END_PARTITION</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ENFORCED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ENUM</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EQUALS</td>
<td>reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ESCAPE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EVENT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EVERY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXCEPT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXCEPTION</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXCLUDE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXCLUDING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXCLUSIVE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EXEC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXECUTE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXISTS</td>
<td>non-reserved</td>
<td>(cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXP</td>
<td></td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXPLAIN</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EXPRESSION</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EXTENSION</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>EXTERNAL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>EXTRACT</td>
<td>non-reserved</td>
<td>(cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FALSE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FAMILY</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FETCH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FILE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>FILTER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>FINAL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>FIRST</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FIRST_VALUE</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FLAG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>FLOAT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FLOOR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>FOLLOWING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>FOR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FORCE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FOREIGN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FORTRAN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>FORWARD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FOUND</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>FRAME_ROW</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FREE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>FREEZE</td>
<td>reserved (can be function or type)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FROM</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>FULL</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>FUNCTION</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>FUNCTIONS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FUSION</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>G</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>GENERAL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>GENERATED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>GET</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>GLOBAL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>GO</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>GOTO</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>GRANT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>GRANTED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>GREATEST</td>
<td>non-reserved (cannot be function or type)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>GROUP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>GROUPING</td>
<td>non-reserved (cannot be function or type)</td>
<td></td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>GROUPS</td>
<td>non-reserved</td>
<td></td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>HANDLER</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>HAVING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>HEADER</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>HEX</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>HIERARCHY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>HOLD</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>HOUR</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ID</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>IDENTITY</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>IF</td>
<td>non-reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IGNORE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>ILIKE</td>
<td>reserved (can be function or type)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>IMMEDIATE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>IMMEDIATELY</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>IMMUTABLE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>IMPLEMENTATION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IMPLICIT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>IMPORT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>IN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INCLUDE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INCLUDING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>INCREMENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>INDENT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INDEX</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INDEXES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INDICATOR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INHERIT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INHERITS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INITIALLY</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INLINE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>INNER</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>INOUT</td>
<td>non-reserved (can not be function or type)</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>INPUT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INSENSITIVE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>INSERT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INSTANCE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>INSTANTIABLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>INSTEAD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>INT</td>
<td>non-reserved (can not be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>-------------</td>
<td>------------</td>
<td>------------</td>
<td>----------</td>
</tr>
<tr>
<td>INTEGER</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INTEGRITY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>INTERSECT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INTERSECTION</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INTERVAL</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INTO</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>INVOKE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>IS</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ISNULL</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ISOLATION</td>
<td>non-reserved (can be function or type)</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>JOIN</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>K</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>KEY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>KEY_MEMBER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>KEY_TYPE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>LABEL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>LAG</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LANGUAGE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LARGE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LAST</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LAST_VALUE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LATERAL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LEAD</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LEADING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LEAKPROOF</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>LEAST</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LEFT</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LENGTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LEVEL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LIBRARY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LIKE</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LIKE_REGEX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LIMIT</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>LINK</td>
<td>non-reserved</td>
<td></td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>LISTEN</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LN</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>LOAD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LOCAL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LOCALTIME</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LOCALTIMESTAMP</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>LOCATION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>LOCATOR</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>LOCK</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LOCKED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LOGGED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LOWER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>M</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MAP</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MAPPING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MATCH</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MATCHED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MATERIALIZED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>MAX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MAXVALUE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MAX_CARDINALITY</td>
<td></td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MEMBER</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MERGE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MESSAGE_LENGTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MESSAGE_OCTET_LENGTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MESSAGE_TEXT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>METHOD</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MIN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MINUTE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MINVALUE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MOD</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MODE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>MODIFIES</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>MODULE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MONTH</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>MORE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>MOVE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>MULTISET</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>MUMPS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>NAMES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NAMESPACE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>NATIONAL</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NATURAL</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NCHAR</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NCLOB</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>NESTING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NEW</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>NEXT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NFC</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NFD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NFKC</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NFKD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NIL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NO</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NONE</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NORMALIZE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>NORMALIZED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>NOT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NOTHING</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>NOTIFY</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>NOTNULL</td>
<td>reserved (can be function or type)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>NOWAIT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>NTH_VALUE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>NTILE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>NULL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NULLABLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>NULLIF</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>NULLS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>NUMBER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
</tbody>
</table>

2318
<table>
<thead>
<tr>
<th>Key Word</th>
<th>PostgreSQL</th>
<th>SQL:2011</th>
<th>SQL:2008</th>
<th>SQL-92</th>
</tr>
</thead>
<tbody>
<tr>
<td>NUMERIC</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OBJECT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OCCURRENCE_REGEX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OCTETS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OCTET_LENGTH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OF</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OFF</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OFFSET</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OIDS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OLD</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ON</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ONLY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OPEN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OPERATOR</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OPTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OPTIONS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ORDER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ORDERING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ORDINALITY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OTHERS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OUT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OUTER</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OUTPUT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OVER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OVERLAPS</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OVERLAY</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>OVERRIDING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OWNED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>OWNER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>P</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>PAD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PARALLEL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PARAMETER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-----------------</td>
<td>------------</td>
<td>------------</td>
<td>---------</td>
</tr>
<tr>
<td>PARAMETER_MODE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARAMETER_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARAMETER_ORDINAL_POSITION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARAMETER_SPECIFIC_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARAMETER_SPECIFIC_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARAMETER_SPECIFIC_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PARSER</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PARTIAL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PARTITION</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>PASCAL</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PASSING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PASSTHROUGH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PASSWORD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PATH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PERCENT</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PERCENTILE_CONT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>PERCENTILE_DISC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>PERCENT_RANK</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>PERIOD</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PERMISSION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PLACING</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PLANS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PLI</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>POLICY</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PORTION</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>POSITION</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>POSITION_REGEX</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>POWER</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>PRECEDES</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>PRECEDING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>PRECISION</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>PREPARE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PREPARED</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PRESERVE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PRIMARY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PRIOR</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PRIVILEGES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PROCEDURAL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PROCEDURE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PROCEDURES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PROGRAM</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PUBLIC</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>PUBLICATION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>QUOTE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RANGE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RANK</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>READ</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>READS</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REAL</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REASSIGN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RECHECK</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RECOVERY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RECURSIVE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REF</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REFERENCES</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REFERENCING</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REFRESH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_AVGX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_AVGY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_COUNT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_INTERCEPT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_R2</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_SLOPE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_SXX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_SXY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REGR_SYY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REINDEX</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RELATIVE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RELEASE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RENAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>REPEATABLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>REPLACE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>REPLICA</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>REQUIRING</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RESET</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>RESPECT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>RESTART</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RESTORE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>RESTRICT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RESULT</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>RETURN</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>RETURNED_CARDINALITY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RETURNED_LENGTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RETURNED_OCTET_LENGTH</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RETURNED_SQLSTATE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RETURNING</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>RETURNS</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>REVOKE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>RIGHT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ROLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROLLBACK</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ROLLUP</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ROUTINE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROUTINES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ROUTINE_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROUTINE_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROUTINE_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROW</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ROWS</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ROW_COUNT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ROW_NUMBER</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>RULE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SAVEPOINT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SCALE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SCHEMAS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCHEMA_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SCOPE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCOPE_CATALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCOPE_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCOPE_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCROLL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SEARCH</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SECOND</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SECTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SECURITY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SELECT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SELECTIVE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SELF</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SENSITIVE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SEQUENCE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SEQUENCES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SERIALIZABLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>SERVER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SERVER_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>SESSION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SESSION_USER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SET</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SETOF</td>
<td>non-reserved</td>
<td>(cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SETS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SHARE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SHOW</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SIMILAR</td>
<td>reserved</td>
<td>reserved</td>
<td>(can be reserved function or type)</td>
<td>reserved</td>
</tr>
<tr>
<td>SIMPLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SIZE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SKIP</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SMALLINT</td>
<td>non-reserved</td>
<td>(cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SNAPSHOT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SOME</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SOURCE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>SPACE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>SPECIFIC</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SPECIFICTYPE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SPECIFIC_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SQL</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SQLCODE</td>
<td></td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SQLERROR</td>
<td></td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SQLEXCEPTION</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SQLSTATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SQLWARNING</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SQRT</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STABLE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>STANDALONE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>START</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>STATE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>STATEMENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>STATIC</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STATISTICS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STDDEV_POP</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STDDEV_SAMP</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STDIN</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>STDOUT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>STORAGE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>STRICT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>STRIP</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>STRUCTURE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>STYLE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>SUB-CLASS_ORIGIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>SUBMULTISET</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SUBSCRIPTION</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SUBSTRING</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SUBSTRING_REGEX</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>SUCCEEDS</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SUM</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SYMMETRIC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SYSID</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SYSTEM</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>SYSTEM_TIME</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SYSTEM_USER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>T</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TABLE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TABLES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>TABLESAMPLE</td>
<td>reserved (can be reserved function or type)</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>TABLESPACE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>TABLE_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>TEMP</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>TEMPLATE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>TEMPORARY</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TEXT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>THEN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TIES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TIME</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TIMESTAMP</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TIME-ZONE_HOUR</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TIME-ZONE_MINUTE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TO</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TOKEN</td>
<td>non-reserved</td>
<td></td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TOP_LEVEL_COUNT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRAILING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSACTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSACTIONS_COMMITTED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSACTIONS_ROLLED_BACK</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSACTION_ACTIVE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSFORM</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSFORMS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSLATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSLATE_REGEX</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRANSLATION</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TREAT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRIGGER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------------------</td>
<td>------------</td>
<td>-----------</td>
<td>-----------</td>
<td>------------</td>
</tr>
<tr>
<td>TRIGGER_CAT ALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>TRIGGER_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>TRIGGER_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>TRIM</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRIM ARRAY</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>TRUE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>TRUNCATE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>TRUSTED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>TYPE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>TYPES</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UESCAPE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNBOUNDED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>UNCOMMITTED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>UNDER</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNENCRIPTED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNION</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>UNIQUE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>UNKNOWN</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>UNLINK</td>
<td>non-reserved</td>
<td></td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>UNLISTEN</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNLOGGED</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNNAMED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>UNNEST</td>
<td>reserved</td>
<td></td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>UNTIL</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>UNTYPED</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>UPDATE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>UPPER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>URI</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>USAGE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>USER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>USER_DEFINED_TYPE_CAT ALOG</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>USER_DEFINED_TYPE_CODE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>USER_DEFINED_TYPE_NAME</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>USER_DEFINED_TYPE_SCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>---------</td>
</tr>
<tr>
<td>USING</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VACUUM</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VALID</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>VALIDATE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VALIDATOR</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VALUE</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VALUES</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VALUE_OF</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VARBINARY</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VARCHAR</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VARIADIC</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VARYING</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VAR_POP</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>VAR_SAMP</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>VERBOSE</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VERSION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>VERSIONING</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VIEW</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>VIEWS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>VOLATILE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>WHEN</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WHENEVER</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WHERE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WHITESPACE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>WIDTH_BUCKET</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>WINDOW</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WITH</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WITHIN</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WITHOUT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WORK</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>WRAPPER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>WRITE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>XML</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLAGG</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>XMLATTRIBUTES</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLBINARY</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>--------------</td>
<td>-------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>XMLCAST</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLCOMMENT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLCONCAT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLDECLARATION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>XMLDOCUMENT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLELEMENT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XML_EXISTS</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLFOREST</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLITERATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLNAMESPACES</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLPARSE</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLPI</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLQUERY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLROOT</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLSCHEMA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>XMLSERIALIZE</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLTABLE</td>
<td>non-reserved (cannot be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLTEXT</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
</tr>
<tr>
<td>XMLVALIDATE</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>YEAR</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>YES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>ZONE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
</tbody>
</table>
Appendix D. SQL Conformance

This section attempts to outline to what extent PostgreSQL conforms to the current SQL standard. The following information is not a full statement of conformance, but it presents the main topics in as much detail as is both reasonable and useful for users.

The formal name of the SQL standard is ISO/IEC 9075 “Database Language SQL”. A revised version of the standard is released from time to time; the most recent update appearing in 2011. The 2011 version is referred to as ISO/IEC 9075:2011, or simply as SQL:2011. The versions prior to that were SQL:2008, SQL:2003, SQL:1999, and SQL-92. Each version replaces the previous one, so claims of conformance to earlier versions have no official merit. PostgreSQL development aims for conformance with the latest official version of the standard where such conformance does not contradict traditional features or common sense. Many of the features required by the SQL standard are supported, though sometimes with slightly differing syntax or function. Further moves towards conformance can be expected over time.

SQL-92 defined three feature sets for conformance: Entry, Intermediate, and Full. Most database management systems claiming SQL standard conformance were conforming at only the Entry level, since the entire set of features in the Intermediate and Full levels was either too voluminous or in conflict with legacy behaviors.

Starting with SQL:1999, the SQL standard defines a large set of individual features rather than the ineffectively broad three levels found in SQL-92. A large subset of these features represents the “Core” features, which every conforming SQL implementation must supply. The rest of the features are purely optional. Some optional features are grouped together to form “packages”, which SQL implementations can claim conformance to, thus claiming conformance to particular groups of features.

The standard versions beginning with SQL:2003 are also split into a number of parts. Each is known by a shorthand name. Note that these parts are not consecutively numbered.

- ISO/IEC 9075-1 Framework (SQL/Framework)
- ISO/IEC 9075-2 Foundation (SQL/Foundation)
- ISO/IEC 9075-3 Call Level Interface (SQL/CLI)
- ISO/IEC 9075-4 Persistent Stored Modules (SQL/PSM)
- ISO/IEC 9075-9 Management of External Data (SQL/MED)
- ISO/IEC 9075-10 Object Language Bindings (SQL/OLB)
- ISO/IEC 9075-11 Information and Definition Schemas (SQL/Schemata)
- ISO/IEC 9075-13 Routines and Types using the Java Language (SQL/JRT)
- ISO/IEC 9075-14 XML-related specifications (SQL/XML)

The PostgreSQL core covers parts 1, 2, 9, 11, and 14. Part 3 is covered by the ODBC driver, and part 13 is covered by the PL/Java plug-in, but exact conformance is currently not being verified for these components. There are currently no implementations of parts 4 and 10 for PostgreSQL.

PostgreSQL supports most of the major features of SQL:2011. Out of 179 mandatory features required for full Core conformance, PostgreSQL conforms to at least 160. In addition, there is a long list of
supported optional features. It might be worth noting that at the time of writing, no current version of any database management system claims full conformance to Core SQL:2011.

In the following two sections, we provide a list of those features that PostgreSQL supports, followed by a list of the features defined in SQL:2011 which are not yet supported in PostgreSQL. Both of these lists are approximate: There might be minor details that are nonconforming for a feature that is listed as supported, and large parts of an unsupported feature might in fact be implemented. The main body of the documentation always contains the most accurate information about what does and does not work.

### Note

Feature codes containing a hyphen are subfeatures. Therefore, if a particular subfeature is not supported, the main feature is listed as unsupported even if some other subfeatures are supported.

## D.1. Supported Features

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Package</th>
<th>Description</th>
<th>Comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>B012</td>
<td>Embedded C</td>
<td></td>
<td></td>
</tr>
<tr>
<td>B021</td>
<td>Direct SQL</td>
<td></td>
<td></td>
</tr>
<tr>
<td>E011</td>
<td>Core</td>
<td>Numeric data types</td>
<td></td>
</tr>
<tr>
<td>E011-01</td>
<td>Core</td>
<td>INTEGER and SMALLINT data types</td>
<td></td>
</tr>
<tr>
<td>E011-02</td>
<td>Core</td>
<td>REAL, DOUBLE PRECISION, and FLOAT data types</td>
<td></td>
</tr>
<tr>
<td>E011-03</td>
<td>Core</td>
<td>DECIMAL and NUMERIC data types</td>
<td></td>
</tr>
<tr>
<td>E011-04</td>
<td>Core</td>
<td>Arithmetic operators</td>
<td></td>
</tr>
<tr>
<td>E011-05</td>
<td>Core</td>
<td>Numeric comparison</td>
<td></td>
</tr>
<tr>
<td>E011-06</td>
<td>Core</td>
<td>Implicit casting among the numeric data types</td>
<td></td>
</tr>
<tr>
<td>E021</td>
<td>Core</td>
<td>Character data types</td>
<td></td>
</tr>
<tr>
<td>E021-01</td>
<td>Core</td>
<td>CHARACTER data type</td>
<td></td>
</tr>
<tr>
<td>E021-02</td>
<td>Core</td>
<td>CHARACTER VARYING data type</td>
<td></td>
</tr>
<tr>
<td>E021-03</td>
<td>Core</td>
<td>Character literals</td>
<td></td>
</tr>
<tr>
<td>E021-04</td>
<td>Core</td>
<td>CHARACTEROPTION_LENGTH function, trims trailing spaces from CHARACTER_LENGTH values before counting</td>
<td></td>
</tr>
<tr>
<td>E021-05</td>
<td>Core</td>
<td>OCTET_LENGTH function</td>
<td></td>
</tr>
<tr>
<td>E021-06</td>
<td>Core</td>
<td>SUBSTRING function</td>
<td></td>
</tr>
<tr>
<td>E021-07</td>
<td>Core</td>
<td>Character concatenation</td>
<td></td>
</tr>
<tr>
<td>E021-08</td>
<td>Core</td>
<td>UPPER and LOWER functions</td>
<td></td>
</tr>
<tr>
<td>E021-09</td>
<td>Core</td>
<td>TRIM function</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>E021-10</td>
<td>Core</td>
<td>Implicit casting among the character string types</td>
<td></td>
</tr>
<tr>
<td>E021-11</td>
<td>Core</td>
<td>POSITION function</td>
<td></td>
</tr>
<tr>
<td>E021-12</td>
<td>Core</td>
<td>Character comparison</td>
<td></td>
</tr>
<tr>
<td>E031</td>
<td>Core</td>
<td>Identifiers</td>
<td></td>
</tr>
<tr>
<td>E031-01</td>
<td>Core</td>
<td>Delimited identifiers</td>
<td></td>
</tr>
<tr>
<td>E031-02</td>
<td>Core</td>
<td>Lower case identifiers</td>
<td></td>
</tr>
<tr>
<td>E031-03</td>
<td>Core</td>
<td>Trailing underscore</td>
<td></td>
</tr>
<tr>
<td>E051</td>
<td>Core</td>
<td>Basic query specification</td>
<td></td>
</tr>
<tr>
<td>E051-01</td>
<td>Core</td>
<td>SELECT DISTINCT</td>
<td></td>
</tr>
<tr>
<td>E051-02</td>
<td>Core</td>
<td>GROUP BY clause</td>
<td></td>
</tr>
<tr>
<td>E051-04</td>
<td>Core</td>
<td>GROUP BY can contain columns not in &lt;select list&gt;</td>
<td></td>
</tr>
<tr>
<td>E051-05</td>
<td>Core</td>
<td>Select list items can be renamed</td>
<td></td>
</tr>
<tr>
<td>E051-06</td>
<td>Core</td>
<td>HAVING clause</td>
<td></td>
</tr>
<tr>
<td>E051-07</td>
<td>Core</td>
<td>Qualified * in select list</td>
<td></td>
</tr>
<tr>
<td>E051-08</td>
<td>Core</td>
<td>Correlation names in the FROM clause</td>
<td></td>
</tr>
<tr>
<td>E051-09</td>
<td>Core</td>
<td>Rename columns in the FROM clause</td>
<td></td>
</tr>
<tr>
<td>E061</td>
<td>Core</td>
<td>Basic predicates and search conditions</td>
<td></td>
</tr>
<tr>
<td>E061-01</td>
<td>Core</td>
<td>Comparison predicate</td>
<td></td>
</tr>
<tr>
<td>E061-02</td>
<td>Core</td>
<td>BETWEEN predicate</td>
<td></td>
</tr>
<tr>
<td>E061-03</td>
<td>Core</td>
<td>IN predicate with list of values</td>
<td></td>
</tr>
<tr>
<td>E061-04</td>
<td>Core</td>
<td>LIKE predicate</td>
<td></td>
</tr>
<tr>
<td>E061-05</td>
<td>Core</td>
<td>LIKE predicate ESCAPE clause</td>
<td></td>
</tr>
<tr>
<td>E061-06</td>
<td>Core</td>
<td>NULL predicate</td>
<td></td>
</tr>
<tr>
<td>E061-07</td>
<td>Core</td>
<td>Quantified comparison predicate</td>
<td></td>
</tr>
<tr>
<td>E061-08</td>
<td>Core</td>
<td>EXISTS predicate</td>
<td></td>
</tr>
<tr>
<td>E061-09</td>
<td>Core</td>
<td>Subqueries in comparison predicate</td>
<td></td>
</tr>
<tr>
<td>E061-11</td>
<td>Core</td>
<td>Subqueries in IN predicate</td>
<td></td>
</tr>
<tr>
<td>E061-12</td>
<td>Core</td>
<td>Subqueries in quantified comparison predicate</td>
<td></td>
</tr>
<tr>
<td>E061-13</td>
<td>Core</td>
<td>Correlated subqueries</td>
<td></td>
</tr>
<tr>
<td>E061-14</td>
<td>Core</td>
<td>Search condition</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>E071</td>
<td>Core</td>
<td>Basic query expressions</td>
<td></td>
</tr>
<tr>
<td>E071-01</td>
<td>Core</td>
<td>UNION DISTINCT table operator</td>
<td></td>
</tr>
<tr>
<td>E071-02</td>
<td>Core</td>
<td>UNION ALL table operator</td>
<td></td>
</tr>
<tr>
<td>E071-03</td>
<td>Core</td>
<td>EXCEPT DISTINCT table operator</td>
<td></td>
</tr>
<tr>
<td>E071-05</td>
<td>Core</td>
<td>Columns combined via table operators need not have exactly the same data type</td>
<td></td>
</tr>
<tr>
<td>E071-06</td>
<td>Core</td>
<td>Table operators in subqueries</td>
<td></td>
</tr>
<tr>
<td>E081</td>
<td>Core</td>
<td>Basic Privileges</td>
<td></td>
</tr>
<tr>
<td>E081-01</td>
<td>Core</td>
<td>SELECT privilege</td>
<td></td>
</tr>
<tr>
<td>E081-02</td>
<td>Core</td>
<td>DELETE privilege</td>
<td></td>
</tr>
<tr>
<td>E081-03</td>
<td>Core</td>
<td>INSERT privilege at the table level</td>
<td></td>
</tr>
<tr>
<td>E081-04</td>
<td>Core</td>
<td>UPDATE privilege at the table level</td>
<td></td>
</tr>
<tr>
<td>E081-05</td>
<td>Core</td>
<td>UPDATE privilege at the column level</td>
<td></td>
</tr>
<tr>
<td>E081-06</td>
<td>Core</td>
<td>REFERENCES privilege at the table level</td>
<td></td>
</tr>
<tr>
<td>E081-07</td>
<td>Core</td>
<td>REFERENCES privilege at the column level</td>
<td></td>
</tr>
<tr>
<td>E081-08</td>
<td>Core</td>
<td>WITH GRANT OPTION</td>
<td></td>
</tr>
<tr>
<td>E081-09</td>
<td>Core</td>
<td>USAGE privilege</td>
<td></td>
</tr>
<tr>
<td>E081-10</td>
<td>Core</td>
<td>EXECUTE privilege</td>
<td></td>
</tr>
<tr>
<td>E091</td>
<td>Core</td>
<td>Set functions</td>
<td></td>
</tr>
<tr>
<td>E091-01</td>
<td>Core</td>
<td>AVG</td>
<td></td>
</tr>
<tr>
<td>E091-02</td>
<td>Core</td>
<td>COUNT</td>
<td></td>
</tr>
<tr>
<td>E091-03</td>
<td>Core</td>
<td>MAX</td>
<td></td>
</tr>
<tr>
<td>E091-04</td>
<td>Core</td>
<td>MIN</td>
<td></td>
</tr>
<tr>
<td>E091-05</td>
<td>Core</td>
<td>SUM</td>
<td></td>
</tr>
<tr>
<td>E091-06</td>
<td>Core</td>
<td>ALL quantifier</td>
<td></td>
</tr>
<tr>
<td>E091-07</td>
<td>Core</td>
<td>DISTINCT quantifier</td>
<td></td>
</tr>
<tr>
<td>E101</td>
<td>Core</td>
<td>Basic data manipulation</td>
<td></td>
</tr>
<tr>
<td>E101-01</td>
<td>Core</td>
<td>INSERT statement</td>
<td></td>
</tr>
<tr>
<td>E101-03</td>
<td>Core</td>
<td>Searched UPDATE statement</td>
<td></td>
</tr>
<tr>
<td>E101-04</td>
<td>Core</td>
<td>Searched DELETE statement</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>E111</td>
<td>Core</td>
<td>Single row SELECT statement</td>
<td></td>
</tr>
<tr>
<td>E121</td>
<td>Core</td>
<td>Basic cursor support</td>
<td></td>
</tr>
<tr>
<td>E121-01</td>
<td>Core</td>
<td>DECLARE CURSOR</td>
<td></td>
</tr>
<tr>
<td>E121-02</td>
<td>Core</td>
<td>ORDER BY columns need not be in select list</td>
<td></td>
</tr>
<tr>
<td>E121-03</td>
<td>Core</td>
<td>Value expressions in ORDER BY clause</td>
<td></td>
</tr>
<tr>
<td>E121-04</td>
<td>Core</td>
<td>OPEN statement</td>
<td></td>
</tr>
<tr>
<td>E121-06</td>
<td>Core</td>
<td>POSITIONED UPDATE statement</td>
<td></td>
</tr>
<tr>
<td>E121-07</td>
<td>Core</td>
<td>POSITIONED DELETE statement</td>
<td></td>
</tr>
<tr>
<td>E121-08</td>
<td>Core</td>
<td>CLOSE statement</td>
<td></td>
</tr>
<tr>
<td>E121-10</td>
<td>Core</td>
<td>FETCH statement implicit NEXT</td>
<td></td>
</tr>
<tr>
<td>E121-17</td>
<td>Core</td>
<td>WITH HOLD cursors</td>
<td></td>
</tr>
<tr>
<td>E131</td>
<td>Core</td>
<td>Null value support (nulls in lieu of values)</td>
<td></td>
</tr>
<tr>
<td>E141</td>
<td>Core</td>
<td>Basic integrity constraints</td>
<td></td>
</tr>
<tr>
<td>E141-01</td>
<td>Core</td>
<td>NOT NULL constraints</td>
<td></td>
</tr>
<tr>
<td>E141-02</td>
<td>Core</td>
<td>UNIQUE constraints of NOT NULL columns</td>
<td></td>
</tr>
<tr>
<td>E141-03</td>
<td>Core</td>
<td>PRIMARY KEY constraints</td>
<td></td>
</tr>
<tr>
<td>E141-04</td>
<td>Core</td>
<td>Basic FOREIGN KEY constraint with the NO ACTION default for both referential delete action and referential update action</td>
<td></td>
</tr>
<tr>
<td>E141-06</td>
<td>Core</td>
<td>CHECK constraints</td>
<td></td>
</tr>
<tr>
<td>E141-07</td>
<td>Core</td>
<td>Column defaults</td>
<td></td>
</tr>
<tr>
<td>E141-08</td>
<td>Core</td>
<td>NOT NULL inferred on PRIMARY KEY</td>
<td></td>
</tr>
<tr>
<td>E141-10</td>
<td>Core</td>
<td>Names in a foreign key can be specified in any order</td>
<td></td>
</tr>
<tr>
<td>E151</td>
<td>Core</td>
<td>Transaction support</td>
<td></td>
</tr>
<tr>
<td>E151-01</td>
<td>Core</td>
<td>COMMIT statement</td>
<td></td>
</tr>
<tr>
<td>E151-02</td>
<td>Core</td>
<td>ROLLBACK statement</td>
<td></td>
</tr>
<tr>
<td>E152</td>
<td>Core</td>
<td>Basic SET TRANSACTION statement</td>
<td></td>
</tr>
<tr>
<td>E152-01</td>
<td>Core</td>
<td>SET TRANSACTION statement: ISOLATION</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>E152-02</td>
<td>Core</td>
<td>LEVEL SERIALIZABLE clause</td>
<td></td>
</tr>
<tr>
<td>E153</td>
<td>Core</td>
<td>SET TRANSACTION statement: READ ONLY and READ WRITE clauses</td>
<td>Updatable queries with subqueries</td>
</tr>
<tr>
<td>E161</td>
<td>Core</td>
<td>SQL comments using leading double minus</td>
<td></td>
</tr>
<tr>
<td>E171</td>
<td>Core</td>
<td>SQLSTATE support</td>
<td></td>
</tr>
<tr>
<td>F021</td>
<td>Core</td>
<td>Basic information schema</td>
<td></td>
</tr>
<tr>
<td>F021-01</td>
<td>Core</td>
<td>COLUMNS view</td>
<td></td>
</tr>
<tr>
<td>F021-02</td>
<td>Core</td>
<td>TABLES view</td>
<td></td>
</tr>
<tr>
<td>F021-03</td>
<td>Core</td>
<td>VIEWS view</td>
<td></td>
</tr>
<tr>
<td>F021-04</td>
<td>Core</td>
<td>TABLE_CONSTRAINTS view</td>
<td></td>
</tr>
<tr>
<td>F021-05</td>
<td>Core</td>
<td>REFERENTIAL_CONSTRAINTS view</td>
<td></td>
</tr>
<tr>
<td>F021-06</td>
<td>Core</td>
<td>CHECK_CONSTRAINTS view</td>
<td></td>
</tr>
<tr>
<td>F031</td>
<td>Core</td>
<td>Basic schema manipulation</td>
<td></td>
</tr>
<tr>
<td>F031-01</td>
<td>Core</td>
<td>CREATE TABLE statement to create persistent base tables</td>
<td></td>
</tr>
<tr>
<td>F031-02</td>
<td>Core</td>
<td>CREATE VIEW statement</td>
<td></td>
</tr>
<tr>
<td>F031-03</td>
<td>Core</td>
<td>GRANT statement</td>
<td></td>
</tr>
<tr>
<td>F031-04</td>
<td>Core</td>
<td>ALTER TABLE statement: ADD COLUMN clause</td>
<td></td>
</tr>
<tr>
<td>F031-13</td>
<td>Core</td>
<td>DROP TABLE statement: RESTRICT clause</td>
<td></td>
</tr>
<tr>
<td>F031-16</td>
<td>Core</td>
<td>DROP VIEW statement: RESTRICT clause</td>
<td></td>
</tr>
<tr>
<td>F031-19</td>
<td>Core</td>
<td>REVOKE statement: RESTRICT clause</td>
<td></td>
</tr>
<tr>
<td>F032</td>
<td></td>
<td>CASCADE drop behavior</td>
<td></td>
</tr>
<tr>
<td>F033</td>
<td></td>
<td>ALTER TABLE statement: DROP COLUMN clause</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F034</td>
<td></td>
<td>Extended REVOKE statement</td>
<td></td>
</tr>
<tr>
<td>F034-01</td>
<td></td>
<td>REVOKE statement performed by other than the owner of a schema object</td>
<td></td>
</tr>
<tr>
<td>F034-02</td>
<td></td>
<td>REVOKE statement: GRANT OPTION FOR clause</td>
<td></td>
</tr>
<tr>
<td>F034-03</td>
<td></td>
<td>REVOKE statement to revoke a privilege that the grantee has WITH GRANT OPTION</td>
<td></td>
</tr>
<tr>
<td>F041</td>
<td>Core</td>
<td>Basic joined table</td>
<td></td>
</tr>
<tr>
<td>F041-01</td>
<td>Core</td>
<td>Inner join (but not necessarily the INNER keyword)</td>
<td></td>
</tr>
<tr>
<td>F041-02</td>
<td>Core</td>
<td>INNER keyword</td>
<td></td>
</tr>
<tr>
<td>F041-03</td>
<td>Core</td>
<td>LEFT OUTER JOIN</td>
<td></td>
</tr>
<tr>
<td>F041-04</td>
<td>Core</td>
<td>RIGHT OUTER JOIN</td>
<td></td>
</tr>
<tr>
<td>F041-05</td>
<td>Core</td>
<td>Outer joins can be nested</td>
<td></td>
</tr>
<tr>
<td>F041-07</td>
<td>Core</td>
<td>The inner table in a left or right outer join can also be used in an inner join</td>
<td></td>
</tr>
<tr>
<td>F041-08</td>
<td>Core</td>
<td>All comparison operators are supported (rather than just =)</td>
<td></td>
</tr>
<tr>
<td>F051</td>
<td>Core</td>
<td>Basic date and time</td>
<td></td>
</tr>
<tr>
<td>F051-01</td>
<td>Core</td>
<td>DATE data type (including support of DATE literal)</td>
<td></td>
</tr>
<tr>
<td>F051-02</td>
<td>Core</td>
<td>TIME data type (including support of TIME literal) with fractional seconds precision of at least 0</td>
<td></td>
</tr>
<tr>
<td>F051-03</td>
<td>Core</td>
<td>TIMESTAMP data type (including support of TIMESTAMP literal) with fractional seconds precision of at least 0 and 6</td>
<td></td>
</tr>
<tr>
<td>F051-04</td>
<td>Core</td>
<td>Comparison predicate on DATE, TIME, and TIMESTAMP data types</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>-----------</td>
<td>-----------------------------------------------------------------------------</td>
<td>--------------------------------------------</td>
</tr>
<tr>
<td>F051-05</td>
<td>Core</td>
<td>Explicit CAST between datetime types and character string types</td>
<td></td>
</tr>
<tr>
<td>F051-06</td>
<td>Core</td>
<td>CURRENT_DATE</td>
<td></td>
</tr>
<tr>
<td>F051-07</td>
<td>Core</td>
<td>LOCALTIME</td>
<td></td>
</tr>
<tr>
<td>F051-08</td>
<td>Core</td>
<td>LOCALTIMESTAMP</td>
<td></td>
</tr>
<tr>
<td>F052</td>
<td>Enhanced</td>
<td>Intervals and datetime arithmetic</td>
<td></td>
</tr>
<tr>
<td>F053</td>
<td></td>
<td>OVERLAPS predicate</td>
<td></td>
</tr>
<tr>
<td>F081</td>
<td>Core</td>
<td>UNION and EXCEPT in views</td>
<td></td>
</tr>
<tr>
<td>F111</td>
<td></td>
<td>Isolation levels other than SERIALIZABLE</td>
<td></td>
</tr>
<tr>
<td>F111-01</td>
<td></td>
<td>READ UNCOMMITTED isolation level</td>
<td></td>
</tr>
<tr>
<td>F111-02</td>
<td></td>
<td>READ COMMITTED isolation level</td>
<td></td>
</tr>
<tr>
<td>F111-03</td>
<td></td>
<td>REPEATABLE READ isolation level</td>
<td></td>
</tr>
<tr>
<td>F131</td>
<td>Core</td>
<td>Grouped operations</td>
<td></td>
</tr>
<tr>
<td>F131-01</td>
<td>Core</td>
<td>WHERE, GROUP BY, and HAVING clauses supported in queries with grouped views</td>
<td></td>
</tr>
<tr>
<td>F131-02</td>
<td>Core</td>
<td>Multiple tables supported in queries with grouped views</td>
<td></td>
</tr>
<tr>
<td>F131-03</td>
<td>Core</td>
<td>Set functions supported in queries with grouped views</td>
<td></td>
</tr>
<tr>
<td>F131-04</td>
<td>Core</td>
<td>Subqueries with GROUP BY and HAVING clauses and grouped views</td>
<td></td>
</tr>
<tr>
<td>F131-05</td>
<td>Core</td>
<td>Single row SELECT with GROUP BY and HAVING clauses and grouped views</td>
<td></td>
</tr>
<tr>
<td>F171</td>
<td></td>
<td>Multiple schemas per user</td>
<td></td>
</tr>
<tr>
<td>F191</td>
<td>Enhanced</td>
<td>Referential delete actions</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>management</td>
<td></td>
<td></td>
</tr>
<tr>
<td>F200</td>
<td></td>
<td>TRUNCATE TABLE statement</td>
<td></td>
</tr>
<tr>
<td>F201</td>
<td>Core</td>
<td>CAST function</td>
<td></td>
</tr>
<tr>
<td>F202</td>
<td></td>
<td>TRUNCATE TABLE: identity column restart option</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F221</td>
<td>Core</td>
<td>Explicit defaults</td>
<td></td>
</tr>
<tr>
<td>F222</td>
<td></td>
<td>INSERT statement: DEFAULT VALUES clause</td>
<td></td>
</tr>
<tr>
<td>F231</td>
<td></td>
<td>Privilege tables</td>
<td></td>
</tr>
<tr>
<td>F231-01</td>
<td></td>
<td>TABLE_PRIVILEGES view</td>
<td></td>
</tr>
<tr>
<td>F231-02</td>
<td></td>
<td>COLUMN_PRIVILEGES view</td>
<td></td>
</tr>
<tr>
<td>F231-03</td>
<td></td>
<td>USAGE_PRIVILEGES view</td>
<td></td>
</tr>
<tr>
<td>F251</td>
<td></td>
<td>Domain support</td>
<td></td>
</tr>
<tr>
<td>F261</td>
<td>Core</td>
<td>CASE expression</td>
<td></td>
</tr>
<tr>
<td>F261-01</td>
<td>Core</td>
<td>Simple CASE</td>
<td></td>
</tr>
<tr>
<td>F261-02</td>
<td>Core</td>
<td>Searched CASE</td>
<td></td>
</tr>
<tr>
<td>F261-03</td>
<td>Core</td>
<td>NULLIF</td>
<td></td>
</tr>
<tr>
<td>F261-04</td>
<td>Core</td>
<td>COALESCE</td>
<td></td>
</tr>
<tr>
<td>F262</td>
<td></td>
<td>Extended CASE expression</td>
<td></td>
</tr>
<tr>
<td>F271</td>
<td></td>
<td>Compound character literals</td>
<td></td>
</tr>
<tr>
<td>F281</td>
<td></td>
<td>LIKE enhancements</td>
<td></td>
</tr>
<tr>
<td>F302</td>
<td></td>
<td>INTERSECT table operator</td>
<td></td>
</tr>
<tr>
<td>F302-01</td>
<td></td>
<td>INTERSECT DISTINCT table operator</td>
<td></td>
</tr>
<tr>
<td>F302-02</td>
<td></td>
<td>INTERSECT ALL table operator</td>
<td></td>
</tr>
<tr>
<td>F304</td>
<td></td>
<td>EXCEPT ALL table operator</td>
<td></td>
</tr>
<tr>
<td>F311-01</td>
<td>Core</td>
<td>CREATE SCHEMA</td>
<td></td>
</tr>
<tr>
<td>F311-02</td>
<td>Core</td>
<td>CREATE TABLE for persistent base tables</td>
<td></td>
</tr>
<tr>
<td>F311-03</td>
<td>Core</td>
<td>CREATE VIEW</td>
<td></td>
</tr>
<tr>
<td>F311-04</td>
<td>Core</td>
<td>CREATE VIEW: WITH CHECK OPTION</td>
<td></td>
</tr>
<tr>
<td>F311-05</td>
<td>Core</td>
<td>GRANT statement</td>
<td></td>
</tr>
<tr>
<td>F321</td>
<td></td>
<td>User authorization</td>
<td></td>
</tr>
<tr>
<td>F361</td>
<td></td>
<td>Subprogram support</td>
<td></td>
</tr>
<tr>
<td>F381</td>
<td></td>
<td>Extended schema manipulation</td>
<td></td>
</tr>
<tr>
<td>F381-01</td>
<td></td>
<td>ALTER TABLE statement: ALTER COLUMN clause</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F381-02</td>
<td></td>
<td>ALTER TABLE statement: ADD CONSTRAINT clause</td>
<td></td>
</tr>
<tr>
<td>F381-03</td>
<td></td>
<td>ALTER TABLE statement: DROP CONSTRAINT clause</td>
<td></td>
</tr>
<tr>
<td>F382</td>
<td></td>
<td>Alter column data type</td>
<td></td>
</tr>
<tr>
<td>F383</td>
<td></td>
<td>Set column not null clause</td>
<td></td>
</tr>
<tr>
<td>F384</td>
<td></td>
<td>Drop identity property clause</td>
<td></td>
</tr>
<tr>
<td>F386</td>
<td></td>
<td>Set identity column generation clause</td>
<td></td>
</tr>
<tr>
<td>F391</td>
<td></td>
<td>Long identifiers</td>
<td></td>
</tr>
<tr>
<td>F392</td>
<td></td>
<td>Unicode escapes in identifiers</td>
<td></td>
</tr>
<tr>
<td>F393</td>
<td></td>
<td>Unicode escapes in literals</td>
<td></td>
</tr>
<tr>
<td>F401</td>
<td></td>
<td>Extended joined table</td>
<td></td>
</tr>
<tr>
<td>F401-01</td>
<td></td>
<td>NATURAL JOIN</td>
<td></td>
</tr>
<tr>
<td>F401-02</td>
<td></td>
<td>FULL OUTER JOIN</td>
<td></td>
</tr>
<tr>
<td>F401-04</td>
<td></td>
<td>CROSS JOIN</td>
<td></td>
</tr>
<tr>
<td>F402</td>
<td></td>
<td>Named column joins for LOBs, arrays, and multisets</td>
<td></td>
</tr>
<tr>
<td>F411</td>
<td>Enhanced datetime facilities</td>
<td>Time zone specification differences regarding literal interpretation</td>
<td></td>
</tr>
<tr>
<td>F421</td>
<td></td>
<td>National character</td>
<td></td>
</tr>
<tr>
<td>F431</td>
<td></td>
<td>Read-only scrollable cursors</td>
<td></td>
</tr>
<tr>
<td>F431-01</td>
<td></td>
<td>FETCH with explicit NEXT</td>
<td></td>
</tr>
<tr>
<td>F431-02</td>
<td></td>
<td>FETCH FIRST</td>
<td></td>
</tr>
<tr>
<td>F431-03</td>
<td></td>
<td>FETCH LAST</td>
<td></td>
</tr>
<tr>
<td>F431-04</td>
<td></td>
<td>FETCH PRIOR</td>
<td></td>
</tr>
<tr>
<td>F431-05</td>
<td></td>
<td>FETCH ABSOLUTE</td>
<td></td>
</tr>
<tr>
<td>F431-06</td>
<td></td>
<td>FETCH RELATIVE</td>
<td></td>
</tr>
<tr>
<td>F441</td>
<td></td>
<td>Extended set function support</td>
<td></td>
</tr>
<tr>
<td>F442</td>
<td></td>
<td>Mixed column references in set functions</td>
<td></td>
</tr>
<tr>
<td>F471</td>
<td>Core</td>
<td>Scalar subquery values</td>
<td></td>
</tr>
<tr>
<td>F481</td>
<td>Core</td>
<td>Expanded NULL predicate</td>
<td></td>
</tr>
<tr>
<td>F491</td>
<td>Enhanced integrity management</td>
<td>Constraint management</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F501</td>
<td>Core</td>
<td>Features and conformance views</td>
<td></td>
</tr>
<tr>
<td>F501-01</td>
<td>Core</td>
<td>SQL_FEATURES view</td>
<td></td>
</tr>
<tr>
<td>F501-02</td>
<td>Core</td>
<td>SQL_SIZING view</td>
<td></td>
</tr>
<tr>
<td>F501-03</td>
<td>Core</td>
<td>SQL_LANGUAGES view</td>
<td></td>
</tr>
<tr>
<td>F502</td>
<td></td>
<td>Enhanced documentation tables</td>
<td></td>
</tr>
<tr>
<td>F502-01</td>
<td></td>
<td>SQL_SIZING_PROFILES view</td>
<td></td>
</tr>
<tr>
<td>F502-02</td>
<td></td>
<td>SQL_IMPLEMENTATION_INFO view</td>
<td></td>
</tr>
<tr>
<td>F502-03</td>
<td></td>
<td>SQL_PACKAGES view</td>
<td></td>
</tr>
<tr>
<td>F531</td>
<td></td>
<td>Temporary tables</td>
<td></td>
</tr>
<tr>
<td>F555</td>
<td>Enhanced datetime facilities</td>
<td>Enhanced seconds precision</td>
<td></td>
</tr>
<tr>
<td>F561</td>
<td></td>
<td>Full value expressions</td>
<td></td>
</tr>
<tr>
<td>F571</td>
<td></td>
<td>Truth value tests</td>
<td></td>
</tr>
<tr>
<td>F591</td>
<td></td>
<td>Derived tables</td>
<td></td>
</tr>
<tr>
<td>F611</td>
<td></td>
<td>Indicator data types</td>
<td></td>
</tr>
<tr>
<td>F641</td>
<td></td>
<td>Row and table constructors</td>
<td></td>
</tr>
<tr>
<td>F651</td>
<td></td>
<td>Catalog name qualifiers</td>
<td></td>
</tr>
<tr>
<td>F661</td>
<td></td>
<td>Simple tables</td>
<td></td>
</tr>
<tr>
<td>F672</td>
<td></td>
<td>Retrospective check constraints</td>
<td></td>
</tr>
<tr>
<td>F690</td>
<td></td>
<td>Collation support but no character set support</td>
<td></td>
</tr>
<tr>
<td>F692</td>
<td></td>
<td>Extended collation support</td>
<td></td>
</tr>
<tr>
<td>F701</td>
<td>Enhanced integrity management</td>
<td>Referential update actions</td>
<td></td>
</tr>
<tr>
<td>F711</td>
<td></td>
<td>ALTER domain</td>
<td></td>
</tr>
<tr>
<td>F731</td>
<td></td>
<td>INSERT column privileges</td>
<td></td>
</tr>
<tr>
<td>F751</td>
<td></td>
<td>View CHECK enhancements</td>
<td></td>
</tr>
<tr>
<td>F761</td>
<td></td>
<td>Session management</td>
<td></td>
</tr>
<tr>
<td>F762</td>
<td></td>
<td>CURRENT_CATALOG</td>
<td></td>
</tr>
<tr>
<td>F763</td>
<td></td>
<td>CURRENT_SCHEMA</td>
<td></td>
</tr>
<tr>
<td>F771</td>
<td></td>
<td>Connection management</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F781</td>
<td></td>
<td>Self-referencing operations</td>
<td></td>
</tr>
<tr>
<td>F791</td>
<td></td>
<td>Insensitive cursors</td>
<td></td>
</tr>
<tr>
<td>F801</td>
<td></td>
<td>Full set function</td>
<td></td>
</tr>
<tr>
<td>F850</td>
<td></td>
<td>Top-level <code>&lt;order by clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F851</td>
<td></td>
<td><code>&lt;order by clause&gt;</code> in subqueries</td>
<td></td>
</tr>
<tr>
<td>F852</td>
<td></td>
<td>Top-level <code>&lt;order by clause&gt;</code> in views</td>
<td></td>
</tr>
<tr>
<td>F855</td>
<td></td>
<td>Nested <code>&lt;order by clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F856</td>
<td></td>
<td>Nested <code>&lt;fetch first clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F857</td>
<td></td>
<td>Top-level <code>&lt;fetch first clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F858</td>
<td></td>
<td><code>&lt;fetch first clause&gt;</code> in subqueries</td>
<td></td>
</tr>
<tr>
<td>F859</td>
<td></td>
<td>Top-level <code>&lt;fetch first clause&gt;</code> in views</td>
<td></td>
</tr>
<tr>
<td>F860</td>
<td></td>
<td><code>&lt;fetch first row count&gt;</code> in <code>&lt;fetch first clause&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F861</td>
<td></td>
<td>Top-level <code>&lt;result offset clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F862</td>
<td></td>
<td><code>&lt;result offset clause&gt;</code> in subqueries</td>
<td></td>
</tr>
<tr>
<td>F863</td>
<td></td>
<td>Nested <code>&lt;result offset clause&gt;</code> in <code>&lt;query expression&gt;</code></td>
<td></td>
</tr>
<tr>
<td>F864</td>
<td></td>
<td>Top-level <code>&lt;result offset clause&gt;</code> in views</td>
<td></td>
</tr>
<tr>
<td>F865</td>
<td></td>
<td><code>&lt;offset row count&gt;</code> in <code>&lt;result offset clause&gt;</code></td>
<td></td>
</tr>
<tr>
<td>S071</td>
<td>Enhanced object support</td>
<td>SQL paths in function and type name resolution</td>
<td></td>
</tr>
<tr>
<td>S092</td>
<td></td>
<td>Arrays of user-defined types</td>
<td></td>
</tr>
<tr>
<td>S095</td>
<td></td>
<td>Array constructors by query</td>
<td></td>
</tr>
<tr>
<td>S096</td>
<td></td>
<td>Optional array bounds</td>
<td></td>
</tr>
<tr>
<td>S098</td>
<td></td>
<td>ARRAY_AGG</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>--------------------------</td>
<td>----------------------------------</td>
<td>-----------------------------------</td>
</tr>
<tr>
<td>S111</td>
<td>Enhanced object support</td>
<td>ONLY in query expressions</td>
<td></td>
</tr>
<tr>
<td>S201</td>
<td>SQL-invoked routines on arrays</td>
<td></td>
<td></td>
</tr>
<tr>
<td>S201-01</td>
<td>Array parameters</td>
<td></td>
<td></td>
</tr>
<tr>
<td>S201-02</td>
<td>Array as result type of functions</td>
<td></td>
<td></td>
</tr>
<tr>
<td>S211</td>
<td>Enhanced object support</td>
<td>User-defined cast functions</td>
<td></td>
</tr>
<tr>
<td>S301</td>
<td>Enhanced UNNEST</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T031</td>
<td>BOOLEAN data type</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T071</td>
<td>BIGINT data type</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T121</td>
<td>WITH (excluding RECURSIVE) in query expression</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T122</td>
<td>WITH (excluding RECURSIVE) in subquery</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T131</td>
<td>Recursive query</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T132</td>
<td>Recursive query in subquery</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T141</td>
<td>SIMILAR predicate</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T151</td>
<td>DISTINCT predicate</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T152</td>
<td>DISTINCT predicate with negation</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T171</td>
<td>LIKE clause in table definition</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T172</td>
<td>AS subquery clause in table definition</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T173</td>
<td>Extended LIKE clause in table definition</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T174</td>
<td>Identity columns</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T177</td>
<td>Sequence generator support: simple restart option</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T178</td>
<td>Identity columns: simple restart option</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T191</td>
<td>Referential action</td>
<td>RESTRICT</td>
<td></td>
</tr>
<tr>
<td>T201</td>
<td>Comparable data types for referential constraints</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T211-01</td>
<td>Triggers activated on UPDATE, INSERT, or DELETE of one base table</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------------------------------</td>
<td>-----------------------------------------------</td>
<td>-------------------------------------------------------------------------</td>
</tr>
<tr>
<td>T211-02</td>
<td>Active database, Enhanced</td>
<td>BEFORE triggers</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity management</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T211-03</td>
<td>Active database, Enhanced</td>
<td>AFTER triggers</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity management</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T211-04</td>
<td>Active database, Enhanced</td>
<td>FOR EACH ROW triggers</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity management</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T211-05</td>
<td>Active database, Enhanced</td>
<td>Ability to specify a search condition that</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity management</td>
<td>must be true before the trigger is invoked</td>
<td></td>
</tr>
<tr>
<td>T211-07</td>
<td>Active database, Enhanced</td>
<td>TRIGGER privilege</td>
<td></td>
</tr>
<tr>
<td></td>
<td>integrity management</td>
<td></td>
<td></td>
</tr>
<tr>
<td>T212</td>
<td>Enhanced integrity management</td>
<td>Enhanced trigger capability</td>
<td></td>
</tr>
<tr>
<td>T213</td>
<td>Enhanced integrity management</td>
<td>INSTEAD OF triggers</td>
<td></td>
</tr>
<tr>
<td>T231</td>
<td>Enhanced integrity management</td>
<td>Sensitive cursors</td>
<td></td>
</tr>
<tr>
<td>T241</td>
<td>Enhanced integrity management</td>
<td>START TRANSACTION statement</td>
<td></td>
</tr>
<tr>
<td>T271</td>
<td>Enhanced integrity management</td>
<td>Savepoints</td>
<td></td>
</tr>
<tr>
<td>T281</td>
<td>Enhanced integrity management</td>
<td>SELECT privilege with column granularity</td>
<td></td>
</tr>
<tr>
<td>T285</td>
<td>Enhanced integrity management</td>
<td>Enhanced derived column names</td>
<td></td>
</tr>
<tr>
<td>T312</td>
<td>Enhanced integrity management</td>
<td>OVERLAY function</td>
<td></td>
</tr>
<tr>
<td>T321-01</td>
<td>Core</td>
<td>User-defined functions with no overloading</td>
<td></td>
</tr>
<tr>
<td>T321-02</td>
<td>Core</td>
<td>User-defined stored procedures with no overloading</td>
<td></td>
</tr>
<tr>
<td>T321-03</td>
<td>Core</td>
<td>Function invocation</td>
<td></td>
</tr>
<tr>
<td>T321-04</td>
<td>Core</td>
<td>CALL statement</td>
<td></td>
</tr>
<tr>
<td>T321-06</td>
<td>Core</td>
<td>ROUTINES view</td>
<td></td>
</tr>
<tr>
<td>T321-07</td>
<td>Core</td>
<td>PARAMETERS view</td>
<td></td>
</tr>
<tr>
<td>T323</td>
<td>Enhanced integrity management</td>
<td>Explicit security for external routines</td>
<td></td>
</tr>
<tr>
<td>T325</td>
<td>Enhanced integrity management</td>
<td>Qualified SQL parameter references</td>
<td></td>
</tr>
<tr>
<td>T331</td>
<td>Enhanced integrity management</td>
<td>Basic roles</td>
<td></td>
</tr>
<tr>
<td>T341</td>
<td>Enhanced integrity management</td>
<td>Overloading of SQL-invoked functions and</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>procedures</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>T351</td>
<td></td>
<td>Bracketed SQL comments (/*...*/ comments)</td>
<td></td>
</tr>
<tr>
<td>T431</td>
<td>OLAP</td>
<td>Extended grouping capabilities</td>
<td></td>
</tr>
<tr>
<td>T432</td>
<td></td>
<td>Nested and concatenated GROUPING SETS</td>
<td></td>
</tr>
<tr>
<td>T433</td>
<td></td>
<td>Multiargument GROUPING function</td>
<td></td>
</tr>
<tr>
<td>T441</td>
<td></td>
<td>ABS and MOD functions</td>
<td></td>
</tr>
<tr>
<td>T461</td>
<td></td>
<td>Symmetric BETWEEN predicate</td>
<td></td>
</tr>
<tr>
<td>T491</td>
<td></td>
<td>LATERAL derived table</td>
<td></td>
</tr>
<tr>
<td>T501</td>
<td></td>
<td>Enhanced EXISTS predicate</td>
<td></td>
</tr>
<tr>
<td>T521</td>
<td></td>
<td>Named arguments in CALL statement</td>
<td></td>
</tr>
<tr>
<td>T551</td>
<td></td>
<td>Optional key words for default syntax</td>
<td></td>
</tr>
<tr>
<td>T581</td>
<td></td>
<td>Regular expression substring function</td>
<td></td>
</tr>
<tr>
<td>T591</td>
<td></td>
<td>UNIQUE constraints of possibly null columns</td>
<td></td>
</tr>
<tr>
<td>T611</td>
<td>OLAP</td>
<td>Elementary OLAP operations</td>
<td></td>
</tr>
<tr>
<td>T613</td>
<td></td>
<td>Sampling</td>
<td></td>
</tr>
<tr>
<td>T614</td>
<td></td>
<td>NTILE function</td>
<td></td>
</tr>
<tr>
<td>T615</td>
<td></td>
<td>LEAD and LAG functions</td>
<td></td>
</tr>
<tr>
<td>T617</td>
<td></td>
<td>FIRST_VALUE and LAST_VALUE function</td>
<td></td>
</tr>
<tr>
<td>T620</td>
<td></td>
<td>WINDOW clause: GROUPS option</td>
<td></td>
</tr>
<tr>
<td>T621</td>
<td></td>
<td>Enhanced numeric functions</td>
<td></td>
</tr>
<tr>
<td>T631</td>
<td>Core</td>
<td>IN predicate with one list element</td>
<td></td>
</tr>
<tr>
<td>T651</td>
<td></td>
<td>SQL-schema statements in SQL routines</td>
<td></td>
</tr>
<tr>
<td>T655</td>
<td></td>
<td>Cyclically dependent routines</td>
<td></td>
</tr>
<tr>
<td>X010</td>
<td></td>
<td>XML type</td>
<td></td>
</tr>
<tr>
<td>X011</td>
<td></td>
<td>Arrays of XML type</td>
<td></td>
</tr>
<tr>
<td>X014</td>
<td></td>
<td>Attributes of XML type</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>X016</td>
<td></td>
<td>Persistent XML values</td>
<td></td>
</tr>
<tr>
<td>X020</td>
<td>XMLConcat</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X031</td>
<td>XMLElement</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X032</td>
<td>XMLForest</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X034</td>
<td>XMLAgg</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X035</td>
<td>XMLAgg: ORDER BY option</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X036</td>
<td>XMLComment</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X037</td>
<td>XMLPI</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X040</td>
<td>Basic table mapping</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X041</td>
<td>Basic table mapping: nulls absent</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X042</td>
<td>Basic table mapping: null as nil</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X043</td>
<td>Basic table mapping: table as forest</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X044</td>
<td>Basic table mapping: table as element</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X045</td>
<td>Basic table mapping: with target namespace</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X046</td>
<td>Basic table mapping: data mapping</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X047</td>
<td>Basic table mapping: metadata mapping</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X048</td>
<td>Basic table mapping: base64 encoding of binary strings</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X049</td>
<td>Basic table mapping: hex encoding of binary strings</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X050</td>
<td>Advanced table mapping</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X051</td>
<td>Advanced table mapping: nulls absent</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X052</td>
<td>Advanced table mapping: null as nil</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X053</td>
<td>Advanced table mapping: table as forest</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X054</td>
<td>Advanced table mapping: table as element</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X055</td>
<td>Advanced table mapping: with target namespace</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X056</td>
<td>Advanced table mapping: data mapping</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X057</td>
<td>Advanced table mapping: metadata mapping</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### D.2. Unsupported Features

The following features defined in SQL:2011 are not implemented in this release of PostgreSQL. In a few cases, equivalent functionality is available.

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Package</th>
<th>Description</th>
<th>Comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>B011</td>
<td></td>
<td>Embedded Ada</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>B013</td>
<td></td>
<td>Embedded COBOL</td>
<td></td>
</tr>
<tr>
<td>B014</td>
<td></td>
<td>Embedded Fortran</td>
<td></td>
</tr>
<tr>
<td>B015</td>
<td></td>
<td>Embedded MUMPS</td>
<td></td>
</tr>
<tr>
<td>B016</td>
<td></td>
<td>Embedded Pascal</td>
<td></td>
</tr>
<tr>
<td>B017</td>
<td></td>
<td>Embedded PL/I</td>
<td></td>
</tr>
<tr>
<td>B031</td>
<td></td>
<td>Basic dynamic SQL</td>
<td></td>
</tr>
<tr>
<td>B032</td>
<td></td>
<td>Extended dynamic SQL</td>
<td></td>
</tr>
<tr>
<td>B032-01</td>
<td></td>
<td>&lt;describe input statement&gt;</td>
<td></td>
</tr>
<tr>
<td>B033</td>
<td></td>
<td>Untyped SQL-invoked function arguments</td>
<td></td>
</tr>
<tr>
<td>B034</td>
<td></td>
<td>Dynamic specification of cursor attributes</td>
<td></td>
</tr>
<tr>
<td>B035</td>
<td></td>
<td>Non-extended descriptor names</td>
<td></td>
</tr>
<tr>
<td>B041</td>
<td></td>
<td>Extensions to embedded SQL exception declarations</td>
<td></td>
</tr>
<tr>
<td>B051</td>
<td></td>
<td>Enhanced execution rights</td>
<td></td>
</tr>
<tr>
<td>B111</td>
<td></td>
<td>Module language Ada</td>
<td></td>
</tr>
<tr>
<td>B112</td>
<td></td>
<td>Module language C</td>
<td></td>
</tr>
<tr>
<td>B113</td>
<td></td>
<td>Module language COBOL</td>
<td></td>
</tr>
<tr>
<td>B114</td>
<td></td>
<td>Module language Fortran</td>
<td></td>
</tr>
<tr>
<td>B115</td>
<td></td>
<td>Module language MUMPS</td>
<td></td>
</tr>
<tr>
<td>B116</td>
<td></td>
<td>Module language Pascal</td>
<td></td>
</tr>
<tr>
<td>B117</td>
<td></td>
<td>Module language PL/I</td>
<td></td>
</tr>
<tr>
<td>B121</td>
<td></td>
<td>Routine language Ada</td>
<td></td>
</tr>
<tr>
<td>B122</td>
<td></td>
<td>Routine language C</td>
<td></td>
</tr>
<tr>
<td>B123</td>
<td></td>
<td>Routine language COBOL</td>
<td></td>
</tr>
<tr>
<td>B124</td>
<td></td>
<td>Routine language Fortran</td>
<td></td>
</tr>
<tr>
<td>B125</td>
<td></td>
<td>Routine language MUMPS</td>
<td></td>
</tr>
<tr>
<td>B126</td>
<td></td>
<td>Routine language Pascal</td>
<td></td>
</tr>
<tr>
<td>B127</td>
<td></td>
<td>Routine language PL/I</td>
<td></td>
</tr>
<tr>
<td>B128</td>
<td></td>
<td>Routine language SQL</td>
<td></td>
</tr>
<tr>
<td>B211</td>
<td></td>
<td>Module language Ada: VARCHAR and NUMERIC support</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td></td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td></td>
</tr>
<tr>
<td>B221</td>
<td></td>
<td>Routine language Ada: VARCHAR and NUMERIC support</td>
<td></td>
</tr>
<tr>
<td>E182</td>
<td>Core</td>
<td>Module language</td>
<td></td>
</tr>
<tr>
<td>F054</td>
<td></td>
<td>TIMESTAMP in DATE type precedence list</td>
<td></td>
</tr>
<tr>
<td>F121</td>
<td></td>
<td>Basic diagnostics management</td>
<td></td>
</tr>
<tr>
<td>F121-01</td>
<td></td>
<td>GET DIAGNOSTICS statement</td>
<td></td>
</tr>
<tr>
<td>F121-02</td>
<td></td>
<td>SET TRANSACTION statement: DIAGNOSTICS SIZE clause</td>
<td></td>
</tr>
<tr>
<td>F122</td>
<td></td>
<td>Enhanced diagnostics management</td>
<td></td>
</tr>
<tr>
<td>F123</td>
<td></td>
<td>All diagnostics</td>
<td></td>
</tr>
<tr>
<td>F181</td>
<td>Core</td>
<td>Multiple module support</td>
<td></td>
</tr>
<tr>
<td>F263</td>
<td></td>
<td>Comma-separated predicates in simple CASE expression</td>
<td></td>
</tr>
<tr>
<td>F291</td>
<td></td>
<td>UNIQUE predicate</td>
<td></td>
</tr>
<tr>
<td>F301</td>
<td></td>
<td>CORRESPONDING in query expressions</td>
<td></td>
</tr>
<tr>
<td>F311</td>
<td>Core</td>
<td>Schema definition statement</td>
<td></td>
</tr>
<tr>
<td>F312</td>
<td></td>
<td>MERGE statement consider INSERT ... ON CONFLICT DO UPDATE</td>
<td></td>
</tr>
<tr>
<td>F313</td>
<td></td>
<td>Enhanced MERGE statement</td>
<td></td>
</tr>
<tr>
<td>F314</td>
<td></td>
<td>MERGE statement with DELETE branch</td>
<td></td>
</tr>
<tr>
<td>F341</td>
<td></td>
<td>Usage tables no ROUTINE_*_USAGE tables</td>
<td></td>
</tr>
<tr>
<td>F385</td>
<td></td>
<td>Drop column generation expression clause</td>
<td></td>
</tr>
<tr>
<td>F394</td>
<td></td>
<td>Optional normal form specification</td>
<td></td>
</tr>
<tr>
<td>F403</td>
<td></td>
<td>Partitioned joined tables</td>
<td></td>
</tr>
<tr>
<td>F451</td>
<td></td>
<td>Character set definition</td>
<td></td>
</tr>
<tr>
<td>F461</td>
<td></td>
<td>Named character sets</td>
<td></td>
</tr>
<tr>
<td>F492</td>
<td></td>
<td>Optional table constraint enforcement</td>
<td></td>
</tr>
<tr>
<td>F521</td>
<td>Enhanced integrity management</td>
<td>Assertions</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>F671</td>
<td>Enhanced integrity management</td>
<td>Subqueries in CHECK</td>
<td>intentionally omitted</td>
</tr>
<tr>
<td>F693</td>
<td></td>
<td>SQL-session and client module collations</td>
<td></td>
</tr>
<tr>
<td>F695</td>
<td></td>
<td>Translation support</td>
<td></td>
</tr>
<tr>
<td>F696</td>
<td></td>
<td>Additional translation documentation</td>
<td></td>
</tr>
<tr>
<td>F721</td>
<td></td>
<td>Deferrable constraints foreign and unique keys only</td>
<td></td>
</tr>
<tr>
<td>F741</td>
<td></td>
<td>Referential MATCH no partial match yet types</td>
<td></td>
</tr>
<tr>
<td>F812</td>
<td>Core</td>
<td>Basic flagging</td>
<td></td>
</tr>
<tr>
<td>F813</td>
<td></td>
<td>Extended flagging</td>
<td></td>
</tr>
<tr>
<td>F821</td>
<td></td>
<td>Local table references</td>
<td></td>
</tr>
<tr>
<td>F831</td>
<td></td>
<td>Full cursor update</td>
<td></td>
</tr>
<tr>
<td>F831-01</td>
<td></td>
<td>Updatable scrollable cursors</td>
<td></td>
</tr>
<tr>
<td>F831-02</td>
<td></td>
<td>Updatable ordered cursors</td>
<td></td>
</tr>
<tr>
<td>F841</td>
<td></td>
<td>LIKE_REGEX predicate</td>
<td></td>
</tr>
<tr>
<td>F842</td>
<td></td>
<td>OCCURRENCES_REGEX function</td>
<td></td>
</tr>
<tr>
<td>F843</td>
<td></td>
<td>POSITION_REGEX function</td>
<td></td>
</tr>
<tr>
<td>F844</td>
<td></td>
<td>SUBSTRING_REGEX function</td>
<td></td>
</tr>
<tr>
<td>F845</td>
<td></td>
<td>TRANSLATE_REGEX function</td>
<td></td>
</tr>
<tr>
<td>F846</td>
<td></td>
<td>Octet support in regular expression operators</td>
<td></td>
</tr>
<tr>
<td>F847</td>
<td></td>
<td>Nonconstant regular expressions</td>
<td></td>
</tr>
<tr>
<td>F866</td>
<td></td>
<td>FETCH FIRST clause: PERCENT option</td>
<td></td>
</tr>
<tr>
<td>F867</td>
<td></td>
<td>FETCH FIRST clause: WITH TIES option</td>
<td></td>
</tr>
<tr>
<td>S011</td>
<td>Core</td>
<td>Distinct data types</td>
<td></td>
</tr>
<tr>
<td>S011-01</td>
<td>Core</td>
<td>USER_DEFINED_TYPES view</td>
<td></td>
</tr>
<tr>
<td>S023</td>
<td>Basic object support</td>
<td>Basic structured types</td>
<td></td>
</tr>
<tr>
<td>S024</td>
<td>Enhanced object support</td>
<td>Enhanced structured types</td>
<td></td>
</tr>
<tr>
<td>S025</td>
<td></td>
<td>Final structured types</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>S026</td>
<td></td>
<td>Self-referencing structured types</td>
<td></td>
</tr>
<tr>
<td>S027</td>
<td></td>
<td>Create method by specific method name</td>
<td></td>
</tr>
<tr>
<td>S028</td>
<td></td>
<td>Permutable UDT options list</td>
<td></td>
</tr>
<tr>
<td>S041</td>
<td>Basic object support</td>
<td>Basic reference types</td>
<td></td>
</tr>
<tr>
<td>S043</td>
<td>Enhanced object support</td>
<td>Enhanced reference types</td>
<td></td>
</tr>
<tr>
<td>S051</td>
<td>Basic object support</td>
<td>Create table of type partially supported</td>
<td></td>
</tr>
<tr>
<td>S081</td>
<td>Enhanced object support</td>
<td>Subtables</td>
<td></td>
</tr>
<tr>
<td>S091</td>
<td></td>
<td>Basic array support partially supported</td>
<td></td>
</tr>
<tr>
<td>S091-01</td>
<td></td>
<td>Arrays of built-in data types</td>
<td></td>
</tr>
<tr>
<td>S091-02</td>
<td></td>
<td>Arrays of distinct types</td>
<td></td>
</tr>
<tr>
<td>S091-03</td>
<td></td>
<td>Array expressions</td>
<td></td>
</tr>
<tr>
<td>S094</td>
<td></td>
<td>Arrays of reference types</td>
<td></td>
</tr>
<tr>
<td>S097</td>
<td></td>
<td>Array element assignment</td>
<td></td>
</tr>
<tr>
<td>S151</td>
<td>Basic object support</td>
<td>Type predicate</td>
<td></td>
</tr>
<tr>
<td>S161</td>
<td>Enhanced object support</td>
<td>Subtype treatment</td>
<td></td>
</tr>
<tr>
<td>S162</td>
<td></td>
<td>Subtype treatment for references</td>
<td></td>
</tr>
<tr>
<td>S202</td>
<td></td>
<td>SQL-invoked routines on multisets</td>
<td></td>
</tr>
<tr>
<td>S231</td>
<td>Enhanced object support</td>
<td>Structured type locators</td>
<td></td>
</tr>
<tr>
<td>S232</td>
<td></td>
<td>Array locators</td>
<td></td>
</tr>
<tr>
<td>S233</td>
<td></td>
<td>Multiset locators</td>
<td></td>
</tr>
<tr>
<td>S241</td>
<td></td>
<td>Transform functions</td>
<td></td>
</tr>
<tr>
<td>S242</td>
<td></td>
<td>Alter transform statement</td>
<td></td>
</tr>
<tr>
<td>S251</td>
<td></td>
<td>User-defined orderings</td>
<td></td>
</tr>
<tr>
<td>S261</td>
<td></td>
<td>Specific type method</td>
<td></td>
</tr>
<tr>
<td>S271</td>
<td></td>
<td>Basic multiset support</td>
<td></td>
</tr>
<tr>
<td>S272</td>
<td></td>
<td>Multisets of user-defined types</td>
<td></td>
</tr>
<tr>
<td>S274</td>
<td></td>
<td>Multisets of reference types</td>
<td></td>
</tr>
<tr>
<td>S275</td>
<td></td>
<td>Advanced multiset support</td>
<td></td>
</tr>
<tr>
<td>S281</td>
<td></td>
<td>Nested collection types</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>S291</td>
<td></td>
<td>Unique constraint on entire row</td>
<td></td>
</tr>
<tr>
<td>S401</td>
<td></td>
<td>Distinct types based on array types</td>
<td></td>
</tr>
<tr>
<td>S402</td>
<td></td>
<td>Distinct types based on distinct types</td>
<td></td>
</tr>
<tr>
<td>S403</td>
<td></td>
<td>AR-RAY_MAX_CARDINALITY</td>
<td></td>
</tr>
<tr>
<td>S404</td>
<td></td>
<td>TRIM_ARRAY</td>
<td></td>
</tr>
<tr>
<td>T011</td>
<td></td>
<td>Timestamp in Information Schema</td>
<td></td>
</tr>
<tr>
<td>T021</td>
<td></td>
<td>BINARY and VARBINARY data types</td>
<td></td>
</tr>
<tr>
<td>T022</td>
<td></td>
<td>Advanced support for BINARY and VARBINARY data types</td>
<td></td>
</tr>
<tr>
<td>T023</td>
<td></td>
<td>Compound binary literal</td>
<td></td>
</tr>
<tr>
<td>T024</td>
<td></td>
<td>Spaces in binary literals</td>
<td></td>
</tr>
<tr>
<td>T041</td>
<td></td>
<td>Basic object support</td>
<td>Basic LOB data type support</td>
</tr>
<tr>
<td>T041-01</td>
<td></td>
<td>Basic object support</td>
<td>BLOB data type</td>
</tr>
<tr>
<td>T041-02</td>
<td></td>
<td>Basic object support</td>
<td>CLOB data type</td>
</tr>
<tr>
<td>T041-03</td>
<td></td>
<td>Basic object support</td>
<td>POSITION, LENGTH, LOWER, TRIM, UPPER, and SUBSTRING functions for LOB data types</td>
</tr>
<tr>
<td>T041-04</td>
<td></td>
<td>Basic object support</td>
<td>Concatenation of LOB data types</td>
</tr>
<tr>
<td>T041-05</td>
<td></td>
<td>Basic object support</td>
<td>LOB locator: non-holdable</td>
</tr>
<tr>
<td>T042</td>
<td></td>
<td>Extended LOB data type support</td>
<td></td>
</tr>
<tr>
<td>T043</td>
<td></td>
<td>Multiplier T</td>
<td></td>
</tr>
<tr>
<td>T044</td>
<td></td>
<td>Multiplier P</td>
<td></td>
</tr>
<tr>
<td>T051</td>
<td></td>
<td>Row types</td>
<td></td>
</tr>
<tr>
<td>T052</td>
<td></td>
<td>MAX and MIN for row types</td>
<td></td>
</tr>
<tr>
<td>T053</td>
<td></td>
<td>Explicit aliases for all-fields reference</td>
<td></td>
</tr>
<tr>
<td>T061</td>
<td></td>
<td>UCS support</td>
<td></td>
</tr>
<tr>
<td>T101</td>
<td></td>
<td>Enhanced nullability determination</td>
<td></td>
</tr>
<tr>
<td>T111</td>
<td></td>
<td>Updatable joins, unions, and columns</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>-----------------------</td>
<td>------------------------------------</td>
<td>-------------------------------------------------------------------------</td>
</tr>
<tr>
<td>T175</td>
<td></td>
<td>Generated columns</td>
<td></td>
</tr>
<tr>
<td>T176</td>
<td></td>
<td>Sequence generator support</td>
<td></td>
</tr>
<tr>
<td>T180</td>
<td></td>
<td>System-versioned tables</td>
<td></td>
</tr>
<tr>
<td>T181</td>
<td></td>
<td>Application-time period tables</td>
<td></td>
</tr>
<tr>
<td>T211</td>
<td>Active database, Enhanced integrity management</td>
<td>Basic trigger capability</td>
<td></td>
</tr>
<tr>
<td>T211-06</td>
<td>Active database, Enhanced integrity management</td>
<td>Support for run-time rules for the interaction of triggers and constraints</td>
<td></td>
</tr>
<tr>
<td>T211-08</td>
<td>Active database, Enhanced integrity management</td>
<td>Multiple triggers for the same event are executed in the order in which they were created in the catalog</td>
<td>intentionally omitted</td>
</tr>
<tr>
<td>T251</td>
<td></td>
<td>SET TRANSACTION statement: LOCAL option</td>
<td></td>
</tr>
<tr>
<td>T261</td>
<td></td>
<td>Chained transactions</td>
<td></td>
</tr>
<tr>
<td>T272</td>
<td></td>
<td>Enhanced savepoint management</td>
<td></td>
</tr>
<tr>
<td>T301</td>
<td></td>
<td>Functional dependencies</td>
<td>partially supported</td>
</tr>
<tr>
<td>T321</td>
<td>Core</td>
<td>Basic SQL-invoked routines</td>
<td></td>
</tr>
<tr>
<td>T321-05</td>
<td>Core</td>
<td>RETURN statement</td>
<td></td>
</tr>
<tr>
<td>T322</td>
<td>PSM</td>
<td>Declared data type attributes</td>
<td></td>
</tr>
<tr>
<td>T324</td>
<td></td>
<td>Explicit security for SQL routines</td>
<td></td>
</tr>
<tr>
<td>T326</td>
<td></td>
<td>Table functions</td>
<td></td>
</tr>
<tr>
<td>T332</td>
<td></td>
<td>Extended roles</td>
<td>mostly supported</td>
</tr>
<tr>
<td>T434</td>
<td></td>
<td>GROUP BY DISTINCT</td>
<td></td>
</tr>
<tr>
<td>T471</td>
<td></td>
<td>Result sets return value</td>
<td></td>
</tr>
<tr>
<td>T472</td>
<td></td>
<td>DESCRIBE CURSOR</td>
<td></td>
</tr>
<tr>
<td>T495</td>
<td></td>
<td>Combined data change different syntax and retrieval</td>
<td></td>
</tr>
<tr>
<td>T502</td>
<td></td>
<td>Period predicates</td>
<td></td>
</tr>
<tr>
<td>T511</td>
<td></td>
<td>Transaction counts</td>
<td></td>
</tr>
<tr>
<td>T522</td>
<td></td>
<td>Default values for IN supported except DE-FAULT key word in invocation</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>T561</td>
<td></td>
<td>Holdable locators</td>
<td></td>
</tr>
<tr>
<td>T571</td>
<td></td>
<td>Array-returning external SQL-invoked functions</td>
<td></td>
</tr>
<tr>
<td>T572</td>
<td></td>
<td>Multiset-returning external SQL-invoked functions</td>
<td></td>
</tr>
<tr>
<td>T601</td>
<td></td>
<td>Local cursor references</td>
<td></td>
</tr>
<tr>
<td>T612</td>
<td></td>
<td>Advanced OLAP operations</td>
<td>some forms supported</td>
</tr>
<tr>
<td>T616</td>
<td></td>
<td>Null treatment option for LEAD and LAG functions</td>
<td></td>
</tr>
<tr>
<td>T618</td>
<td></td>
<td>NTH_VALUE function exists, but some options missing</td>
<td></td>
</tr>
<tr>
<td>T619</td>
<td></td>
<td>Nested window functions</td>
<td></td>
</tr>
<tr>
<td>T641</td>
<td></td>
<td>Multiple column assignment</td>
<td>only some syntax variants supported</td>
</tr>
<tr>
<td>T652</td>
<td></td>
<td>SQL-dynamic statements in SQL routines</td>
<td></td>
</tr>
<tr>
<td>T653</td>
<td></td>
<td>SQL-schema statements in external routines</td>
<td></td>
</tr>
<tr>
<td>T654</td>
<td></td>
<td>SQL-dynamic statements in external routines</td>
<td></td>
</tr>
<tr>
<td>M001</td>
<td></td>
<td>Datalinks</td>
<td></td>
</tr>
<tr>
<td>M002</td>
<td></td>
<td>Datalinks via SQL/CLI</td>
<td></td>
</tr>
<tr>
<td>M003</td>
<td></td>
<td>Datalinks via Embedded SQL</td>
<td></td>
</tr>
<tr>
<td>M004</td>
<td></td>
<td>Foreign data support partially supported</td>
<td></td>
</tr>
<tr>
<td>M005</td>
<td></td>
<td>Foreign schema support</td>
<td></td>
</tr>
<tr>
<td>M006</td>
<td></td>
<td>GetSQLString routine</td>
<td></td>
</tr>
<tr>
<td>M007</td>
<td></td>
<td>TransmitRequest</td>
<td></td>
</tr>
<tr>
<td>M009</td>
<td></td>
<td>GetOpts and GetStatistics routines</td>
<td></td>
</tr>
<tr>
<td>M010</td>
<td></td>
<td>Foreign data wrapper different API support</td>
<td></td>
</tr>
<tr>
<td>M011</td>
<td></td>
<td>Datalinks via Ada</td>
<td></td>
</tr>
<tr>
<td>M012</td>
<td></td>
<td>Datalinks via C</td>
<td></td>
</tr>
<tr>
<td>M013</td>
<td></td>
<td>Datalinks via COBOL</td>
<td></td>
</tr>
<tr>
<td>M014</td>
<td></td>
<td>Datalinks via Fortran</td>
<td></td>
</tr>
<tr>
<td>M015</td>
<td></td>
<td>Datalinks via M</td>
<td></td>
</tr>
<tr>
<td>M016</td>
<td></td>
<td>Datalinks via Pascal</td>
<td></td>
</tr>
<tr>
<td>M017</td>
<td></td>
<td>Datalinks via PL/I</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>M018</td>
<td></td>
<td>Foreign data wrapper interface routines in Ada</td>
<td></td>
</tr>
<tr>
<td>M019</td>
<td></td>
<td>Foreign data wrapper different API interface routines in C</td>
<td></td>
</tr>
<tr>
<td>M020</td>
<td></td>
<td>Foreign data wrapper interface routines in COBOL</td>
<td></td>
</tr>
<tr>
<td>M021</td>
<td></td>
<td>Foreign data wrapper interface routines in Fortran</td>
<td></td>
</tr>
<tr>
<td>M022</td>
<td></td>
<td>Foreign data wrapper interface routines in MUMPS</td>
<td></td>
</tr>
<tr>
<td>M023</td>
<td></td>
<td>Foreign data wrapper interface routines in Pascal</td>
<td></td>
</tr>
<tr>
<td>M024</td>
<td></td>
<td>Foreign data wrapper interface routines in PL/I</td>
<td></td>
</tr>
<tr>
<td>M030</td>
<td></td>
<td>SQL-server foreign data support</td>
<td></td>
</tr>
<tr>
<td>M031</td>
<td></td>
<td>Foreign data wrapper general routines</td>
<td></td>
</tr>
<tr>
<td>X012</td>
<td></td>
<td>Multisets of XML type</td>
<td></td>
</tr>
<tr>
<td>X013</td>
<td></td>
<td>Distinct types of XML type</td>
<td></td>
</tr>
<tr>
<td>X015</td>
<td></td>
<td>Fields of XML type</td>
<td></td>
</tr>
<tr>
<td>X025</td>
<td></td>
<td>XMLCast</td>
<td></td>
</tr>
<tr>
<td>X030</td>
<td></td>
<td>XMLDocument</td>
<td></td>
</tr>
<tr>
<td>X038</td>
<td></td>
<td>XMLText</td>
<td></td>
</tr>
<tr>
<td>X065</td>
<td></td>
<td>XMLParse: BLOB input and CONTENT option</td>
<td></td>
</tr>
<tr>
<td>X066</td>
<td></td>
<td>XMLParse: BLOB input and DOCUMENT option</td>
<td></td>
</tr>
<tr>
<td>X068</td>
<td></td>
<td>XMLSerialize: BOM</td>
<td></td>
</tr>
<tr>
<td>X069</td>
<td></td>
<td>XMLSerialize: INDENT</td>
<td></td>
</tr>
<tr>
<td>X073</td>
<td></td>
<td>XMLSerialize: BLOB serialization and CONTENT option</td>
<td></td>
</tr>
<tr>
<td>X074</td>
<td></td>
<td>XMLSerialize: BLOB serialization and DOCUMENT option</td>
<td></td>
</tr>
<tr>
<td>X075</td>
<td></td>
<td>XMLSerialize: BLOB serialization</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>-----------</td>
<td>------------------------------------</td>
<td>-----------</td>
</tr>
<tr>
<td>X076</td>
<td></td>
<td>XMLSerialize: VERSION</td>
<td></td>
</tr>
<tr>
<td>X077</td>
<td></td>
<td>XMLSerialize: explicit ENCODING option</td>
<td></td>
</tr>
<tr>
<td>X078</td>
<td></td>
<td>XMLSerialize: explicit XML declaration</td>
<td></td>
</tr>
<tr>
<td>X080</td>
<td></td>
<td>Namespaces in XML publishing</td>
<td></td>
</tr>
<tr>
<td>X081</td>
<td></td>
<td>Query-level XML namespace declarations</td>
<td></td>
</tr>
<tr>
<td>X082</td>
<td></td>
<td>XML namespace declarations in DML</td>
<td></td>
</tr>
<tr>
<td>X083</td>
<td></td>
<td>XML namespace declarations in DDL</td>
<td></td>
</tr>
<tr>
<td>X084</td>
<td></td>
<td>XML namespace declarations in compound statements</td>
<td></td>
</tr>
<tr>
<td>X085</td>
<td></td>
<td>Predefined namespace prefixes</td>
<td></td>
</tr>
<tr>
<td>X086</td>
<td></td>
<td>XML namespace declarations in XMLTable</td>
<td></td>
</tr>
<tr>
<td>X091</td>
<td></td>
<td>XML content predicate</td>
<td></td>
</tr>
<tr>
<td>X096</td>
<td></td>
<td>XMLExists XPath only</td>
<td></td>
</tr>
<tr>
<td>X100</td>
<td></td>
<td>Host language support for XML: CONTENT option</td>
<td></td>
</tr>
<tr>
<td>X101</td>
<td></td>
<td>Host language support for XML: DOCUMENT option</td>
<td></td>
</tr>
<tr>
<td>X110</td>
<td></td>
<td>Host language support for XML: VARCHAR mapping</td>
<td></td>
</tr>
<tr>
<td>X111</td>
<td></td>
<td>Host language support for XML: CLOB mapping</td>
<td></td>
</tr>
<tr>
<td>X112</td>
<td></td>
<td>Host language support for XML: BLOB mapping</td>
<td></td>
</tr>
<tr>
<td>X113</td>
<td></td>
<td>Host language support for XML: STRIP WHITESPACE option</td>
<td></td>
</tr>
<tr>
<td>X114</td>
<td></td>
<td>Host language support for XML: PRESERVE WHITESPACE option</td>
<td></td>
</tr>
<tr>
<td>X131</td>
<td></td>
<td>Query-level XMLBINARY clause</td>
<td></td>
</tr>
<tr>
<td>X132</td>
<td></td>
<td>XMLBINARY clause in DML</td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>X133</td>
<td></td>
<td>XMLBINARY clause in DDL</td>
<td></td>
</tr>
<tr>
<td>X134</td>
<td></td>
<td>XMLBINARY clause in compound statements</td>
<td></td>
</tr>
<tr>
<td>X135</td>
<td></td>
<td>XMLBINARY clause in subqueries</td>
<td></td>
</tr>
<tr>
<td>X141</td>
<td></td>
<td>IS VALID predicate: data-driven case</td>
<td></td>
</tr>
<tr>
<td>X142</td>
<td></td>
<td>IS VALID predicate: ACCORDING TO clause</td>
<td></td>
</tr>
<tr>
<td>X143</td>
<td></td>
<td>IS VALID predicate: ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X144</td>
<td></td>
<td>IS VALID predicate: schema location</td>
<td></td>
</tr>
<tr>
<td>X145</td>
<td></td>
<td>IS VALID predicate outside check constraints</td>
<td></td>
</tr>
<tr>
<td>X151</td>
<td></td>
<td>IS VALID predicate with DOCUMENT option</td>
<td></td>
</tr>
<tr>
<td>X152</td>
<td></td>
<td>IS VALID predicate with CONTENT option</td>
<td></td>
</tr>
<tr>
<td>X153</td>
<td></td>
<td>IS VALID predicate with SEQUENCE option</td>
<td></td>
</tr>
<tr>
<td>X155</td>
<td></td>
<td>IS VALID predicate: NAMESPACE without ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X157</td>
<td></td>
<td>IS VALID predicate: NO NAMESPACE with ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X160</td>
<td>Basic Information Schema for registered XML Schemas</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X161</td>
<td>Advanced Information Schema for registered XML Schemas</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X170</td>
<td>XML null handling options</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X171</td>
<td>NIL ON NO CONTENT option</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X181</td>
<td>XML(DOCUMENT(T(UNTYPED))) type</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X182</td>
<td>XML(DOCUMENT(T(ANY))) type</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X190</td>
<td>XML(SEQUENCE) type</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>X191</td>
<td></td>
<td>XML(DOCUMENT-</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>T(XMLSCHEMA))</td>
<td>type</td>
</tr>
<tr>
<td>X192</td>
<td></td>
<td>XML(CONTENT(XM-</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>LSCHEMA)) type</td>
<td></td>
</tr>
<tr>
<td>X200</td>
<td></td>
<td>XMLQuery</td>
<td></td>
</tr>
<tr>
<td>X201</td>
<td></td>
<td>XMLQuery: RETURNING CONTENT</td>
<td></td>
</tr>
<tr>
<td>X202</td>
<td></td>
<td>XMLQuery: RETURNING SE-</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>QUENCE</td>
<td></td>
</tr>
<tr>
<td>X203</td>
<td></td>
<td>XMLQuery: passing a</td>
<td>context item</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X204</td>
<td></td>
<td>XMLQuery: initializing an XQuery</td>
<td>variable</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X205</td>
<td></td>
<td>XMLQuery: EMPTY ON EMPTY</td>
<td>option</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X206</td>
<td></td>
<td>XMLQuery: NULL ON EMPTY</td>
<td>option</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X211</td>
<td></td>
<td>XML 1.1 support</td>
<td></td>
</tr>
<tr>
<td>X221</td>
<td></td>
<td>XML passing mechanism BY VALUE</td>
<td></td>
</tr>
<tr>
<td>X231</td>
<td></td>
<td>XML(CONTENT(UN-</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>TYPED)) type</td>
<td></td>
</tr>
<tr>
<td>X232</td>
<td></td>
<td>XML(CONTENT(ANY)) type</td>
<td></td>
</tr>
<tr>
<td>X241</td>
<td></td>
<td>RETURNING CON-</td>
<td>TENT in XML publishing</td>
</tr>
<tr>
<td></td>
<td></td>
<td>TENT</td>
<td></td>
</tr>
<tr>
<td>X242</td>
<td></td>
<td>RETURNING SE-</td>
<td>QUENCE in XML publishing</td>
</tr>
<tr>
<td></td>
<td></td>
<td>QUENCE</td>
<td></td>
</tr>
<tr>
<td>X251</td>
<td></td>
<td>Persistent XML values of XML(DOCUMENT-</td>
<td>T(UNTYPED)) type</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X252</td>
<td></td>
<td>Persistent XML values of XML(DOCUMENT-</td>
<td>T(ANY)) type</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X253</td>
<td></td>
<td>Persistent XML values of XML(CONTENT(UN-</td>
<td>TYPED)) type</td>
</tr>
<tr>
<td></td>
<td></td>
<td>XML</td>
<td></td>
</tr>
<tr>
<td>X254</td>
<td></td>
<td>Persistent XML values of XML(CONTENT(ANY)) type</td>
<td></td>
</tr>
<tr>
<td>X255</td>
<td></td>
<td>Persistent XML values of XML(SEQUENCE)</td>
<td>type</td>
</tr>
<tr>
<td>Identifier</td>
<td>Package</td>
<td>Description</td>
<td>Comment</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>-------------</td>
<td>---------</td>
</tr>
<tr>
<td>X256</td>
<td></td>
<td>Persistent XML values of XML(DOCUMENT(XMLSCHEMA)) type</td>
<td></td>
</tr>
<tr>
<td>X257</td>
<td></td>
<td>Persistent XML values of XML(CONTENT(XMLSCHEMA)) type</td>
<td></td>
</tr>
<tr>
<td>X260</td>
<td></td>
<td>XML type: ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X261</td>
<td></td>
<td>XML type: NAMESPACE without ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X263</td>
<td></td>
<td>XML type: NO NAMESPACE with ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X264</td>
<td></td>
<td>XML type: schema location</td>
<td></td>
</tr>
<tr>
<td>X271</td>
<td></td>
<td>XMLValidate: data-driven case</td>
<td></td>
</tr>
<tr>
<td>X272</td>
<td></td>
<td>XMLValidate: ACCORDING TO clause</td>
<td></td>
</tr>
<tr>
<td>X273</td>
<td></td>
<td>XMLValidate: ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X274</td>
<td></td>
<td>XMLValidate: schema location</td>
<td></td>
</tr>
<tr>
<td>X281</td>
<td></td>
<td>XMLValidate: DOCUMENT option</td>
<td></td>
</tr>
<tr>
<td>X282</td>
<td></td>
<td>XMLValidate: CONTENT option</td>
<td></td>
</tr>
<tr>
<td>X283</td>
<td></td>
<td>XMLValidate: SEQUENCE option</td>
<td></td>
</tr>
<tr>
<td>X284</td>
<td></td>
<td>XMLValidate: NAMESPACE without ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X286</td>
<td></td>
<td>XMLValidate: NO NAMESPACE with ELEMENT clause</td>
<td></td>
</tr>
<tr>
<td>X300</td>
<td></td>
<td>XMLTable XPath only</td>
<td></td>
</tr>
<tr>
<td>X305</td>
<td></td>
<td>XMLTable: initializing an XQuery variable</td>
<td></td>
</tr>
</tbody>
</table>
Appendix E. Release Notes

The release notes contain the significant changes in each PostgreSQL release, with major features and migration issues listed at the top. The release notes do not contain changes that affect only a few users or changes that are internal and therefore not user-visible. For example, the optimizer is improved in almost every release, but the improvements are usually observed by users as simply faster queries.

A complete list of changes for each release can be obtained by viewing the Git logs for each release. The postgres-committers email list\(^1\) records all source code changes as well. There is also a web interface\(^2\) that shows changes to specific files.

The name appearing next to each item represents the major developer for that item. Of course all changes involve community discussion and patch review, so each item is truly a community effort.

E.1. Release 11.1

**Release date:** 2018-11-08

This release contains a variety of fixes from 11.0. For information about new features in major release 11, see Section E.2.

E.1.1. Migration to Version 11.1

A dump/restore is not required for those running 11.X.

However, if you use the `pg_stat_statements` extension, see the changelog entry below about that.

E.1.2. Changes

- Ensure proper quoting of transition table names when `pg_dump` emits `CREATE TRIGGER ... REFERENCING` commands (Tom Lane)
  
  This oversight could be exploited by an unprivileged user to gain superuser privileges during the next dump/reload or `pg_upgrade` run. (CVE-2018-16850)
- Apply the tablespace specified for a partitioned index when creating a child index (Álvaro Herrera)
  
  Previously, child indexes were always created in the default tablespace.
- Fix NULL handling in parallel hashed multi-batch left joins (Andrew Gierth, Thomas Munro)
  
  Outer-relation rows with null values of the hash key were omitted from the join result.
- Fix incorrect processing of an array-type coercion expression appearing within a `CASE` clause that has a constant test expression (Tom Lane)
- Fix incorrect expansion of tuples lacking recently-added columns (Andrew Dunstan, Amit Langote)
  
  This is known to lead to crashes in triggers on tables with recently-added columns, and could have other symptoms as well.
- Fix bugs with named or defaulted arguments in `CALL` argument lists (Tom Lane, Pavel Stehule)
- Fix strictness check for strict aggregates with `ORDER BY` columns (Andrew Gierth, Andres Freund)

\(^1\) [https://www.postgresql.org/list/postgres-committers/](https://www.postgresql.org/list/postgres-committers/)
\(^2\) [https://git.postgresql.org/gitweb/?p=postgresql.git;a=summary](https://git.postgresql.org/gitweb/?p=postgresql.git;a=summary)
The strictness logic incorrectly ignored rows for which the ORDER BY value(s) were null.

- Disable recheck_on_update optimization (Tom Lane)

  This new-in-v11 feature turns out not to have been ready for prime time. Disable it until something can be done about it.

- Prevent creation of a partition in a trigger attached to its parent table (Amit Langote)

  Ideally we'd allow that, but for the moment it has to be blocked to avoid crashes.

- Fix problems with applying ON COMMIT DELETE ROWS to a partitioned temporary table (Amit Langote)

- Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)

  This bug affected full-text-search operations, as well as contrib/ltrie and contrib/pg_trgm.

- Ensure that the server will process already-received NOTIFY and SIGTERM interrupts before waiting for client input (Jeff Janes, Tom Lane)

- Fix memory leak in repeated SP-GiST index scans (Tom Lane)

  This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- Prevent starting the server with wal_level set to too low a value to support an existing replication slot (Andres Freund)

- Fix psql, as well as documentation examples, to call PQconsumeInput() before each PQnotifies() call (Tom Lane)

  This fixes cases in which psql would not report receipt of a NOTIFY message until after the next command.

- Fix pg_verify_checksums's determination of which files to check the checksums of (Michael Paquier)

  In some cases it complained about files that are not expected to have checksums.

- In contrib/pg_stat_statements, disallow the pg_read_all_stats role from executing pg_stat_statements_reset() (Haribabu Kommi)

  pg_read_all_stats is only meant to grant permission to read statistics, not to change them, so this grant was incorrect.

  To cause this change to take effect, run ALTER EXTENSION pg_stat_statements UPDATE in each database where pg_stat_statements has been installed. (A database freshly created in 11.0 should not need this, but a database upgraded from a previous release probably still contains the old version of pg_stat_statements. The UPDATE command is harmless if the module was already updated.)

- Rename red-black tree support functions to use rbt prefix not rb prefix (Tom Lane)

  This avoids name collisions with Ruby functions, which broke PL/Ruby. It's hoped that there are no other affected extensions.

- Fix build problems on macOS 10.14 (Mojave) (Tom Lane)
Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.

It is now recommended that Perl-related extensions write $(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.

Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.

- Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)

Perl no longer includes the current directory in its search path by default; work around that.

- On Windows, allow the regression tests to be run by an Administrator account (Andrew Dunstan)

To do this safely, pg_regress now gives up any such privileges at startup.

- Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.2. Release 11

E.2.1. Overview

Major enhancements in PostgreSQL 11 include:

- Improvements to partitioning functionality, including:
  - Add support for partitioning by a hash key
  - Add support for PRIMARY KEY, FOREIGN KEY, indexes, and triggers on partitioned tables
  - Allow creation of a “default” partition for storing data that does not match any of the remaining partitions
  - UPDATE statements that change a partition key column now cause affected rows to be moved to the appropriate partitions
  - Improve SELECT performance through enhanced partition elimination strategies during query planning and execution

- Improvements to parallelism, including:
  - CREATE INDEX can now use parallel processing while building a B-tree index
  - Parallelization is now possible in CREATE TABLE ... AS, CREATE MATERIALIZED VIEW, and certain queries using UNION
  - Parallelized hash joins and parallelized sequential scans now perform better
  - SQL stored procedures that support embedded transactions
  - Optional Just-in-Time (JIT) compilation for some SQL code, speeding evaluation of expressions
  - Window functions now support all framing options shown in the SQL:2011 standard, including RANGE distance PRECEDING/FOLLOWING, GROUPS mode, and frame exclusion options
• Covering indexes can now be created, using the INCLUDE clause of CREATE INDEX

• Many other useful performance improvements, including the ability to avoid a table rewrite for ALTER TABLE ... ADD COLUMN with a non-null column default

The above items are explained in more detail in the sections below.

E.2.2. Migration to Version 11

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 11 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

• Make pg_dump dump the properties of a database, not just its contents (Haribabu Kommi)

Previously, attributes of the database itself, such as database-level GRANT/REVOKE permissions and ALTER DATABASE SET variable settings, were only dumped by pg_dumpall. Now pg_dump --create and pg_restore --create will restore these database properties in addition to the objects within the database. pg_dumpall -g now only dumps role- and tablespace-related attributes. pg_dumpall's complete output (without -g) is unchanged.

pg_dump and pg_restore, without --create, no longer dump/restore database-level comments and security labels; those are now treated as properties of the database.

pg_dumpall's output script will now always create databases with their original locale and encoding, and hence will fail if the locale or encoding name is unknown to the destination system. Previously, CREATE DATABASE would be emitted without these specifications if the database locale and encoding matched the old cluster's defaults.

pg_dumpall --clean now restores the original locale and encoding settings of the postgres and template1 databases, as well as those of user-created databases.

• Consider syntactic form when disambiguating function versus column references (Tom Lane)

When \( x \) is a table name or composite column, PostgreSQL has traditionally considered the syntactic forms \( f(x) \) and \( x.f \) to be equivalent, allowing tricks such as writing a function and then using it as though it were a computed-on-demand column. However, if both interpretations are feasible, the column interpretation was always chosen, leading to surprising results if the user intended the function interpretation. Now, if there is ambiguity, the interpretation that matches the syntactic form is chosen.

• Fully enforce uniqueness of table and domain constraint names (Tom Lane)

PostgreSQL expects the names of a table's constraints to be distinct, and likewise for the names of a domain's constraints. However, there was not rigid enforcement of this, and previously there were corner cases where duplicate names could be created.

• Make \( \text{power(numeric, numeric)} \) and \( \text{power(float8, float8)} \) handle NaN inputs according to the POSIX standard (Tom Lane, Dang Minh Huong)

POSIX says that \( \text{NaN}^0 = 1 \) and \( 1^\text{NaN} = 1 \), but all other cases with \( \text{NaN} \) input(s) should return \( \text{NaN} \). \( \text{power(numeric, numeric)} \) just returned \( \text{NaN} \) in all such cases; now it honors the two exceptions. \( \text{power(float8, float8)} \) followed the standard if the C library does; but on some old Unix platforms the library doesn't, and there were also problems on some versions of Windows.

• Prevent \( \text{to_number()} \) from consuming characters when the template separator does not match (Oliver Ford)
Specifically, `SELECT to_number('1234', '9,999')` used to return 134. It will now return 1234. `L` and `TH` now only consume characters that are not digits, positive/negative signs, decimal points, or commas.

- **Fix to_date(), to_number(), and to_timestamp() to skip a character for each template character (Tom Lane)**

  Previously, they skipped one byte for each byte of template character, resulting in strange behavior if either string contained multibyte characters.

- **Adjust the handling of backslashes inside double-quotes in template strings for to_char(), to_number(), and to_timestamp().**

  Such a backslash now escapes the character after it, particularly a double-quote or another backslash.

- **Correctly handle relative path expressions in xmltable(), xpath(), and other XML-handling functions (Markus Winand)**

  Per the SQL standard, relative paths start from the document node of the XML input document, not the root node as these functions previously did.

- **In the extended query protocol, make statement_timeout apply to each Execute message separately, not to all commands before Sync (Tatsuo Ishii, Andres Freund)**

- **Remove the relhaspkey column from system catalog pg_class (Peter Eisentraut)**

  Applications needing to check for a primary key should consult `pg_index`.

- **Replace system catalog pg_proc's proisagg and proiswindow columns with prokind (Peter Eisentraut)**

  This new column more clearly distinguishes functions, procedures, aggregates, and window functions.

- **Correct information schema column tables.table_type to return FOREIGN instead of FOREIGN TABLE (Peter Eisentraut)**

  This new output matches the SQL standard.

- **Change the ps process display labels for background workers to match the pg_stat_activity.backend_type labels (Peter Eisentraut)**

- **Cause large object permission checks to happen during large object open, lo_open(), not when a read or write is attempted (Tom Lane, Michael Paquier)**

  If write access is requested and not available, an error will now be thrown even if the large object is never written to.

- **Prevent non-superusers from reindexing shared catalogs (Michael Paquier, Robert Haas)**

  Previously, database owners were also allowed to do this, but now it is considered outside the bounds of their privileges.

- **Remove deprecated adminpack functions pg_file_read(), pg_file_length(), and pg_logfile_rotate() (Stephen Frost)**

  Equivalent functionality is now present in the core backend. Existing adminpack installs will continue to have access to these functions until they are updated via `ALTER EXTENSION ... UPDATE`.

- **Honor the capitalization of double-quoted command options (Daniel Gustafsson)**
Previously, option names in certain SQL commands were forcibly lower-cased even if entered with double quotes; thus for example "FillFactor" would be accepted as an index storage option, though properly its name is lower-case. Such cases will now generate an error.

- Remove server parameter replacement_sort_tuples (Peter Geoghegan)
  Replacement sorts were determined to be no longer useful.

- Remove WITH clause in CREATE FUNCTION (Michael Paquier)
  PostgreSQL has long supported a more standard-compliant syntax for this capability.

E.2.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 11 and the previous major release.

E.2.3.1. Server

E.2.3.1.1. Partitioning

- Allow the creation of partitions based on hashing a key column (Amul Sul)

- Support indexes on partitioned tables (Álvaro Herrera, Amit Langote)
  An “index” on a partitioned table is not a physical index across the whole partitioned table, but rather a template for automatically creating similar indexes on each partition of the table.

  If the partition key is part of the index's column set, a partitioned index may be declared UNIQUE. It will represent a valid uniqueness constraint across the whole partitioned table, even though each physical index only enforces uniqueness within its own partition.

  The new command ALTER INDEX ATTACH PARTITION causes an existing index on a partition to be associated with a matching index template for its partitioned table. This provides flexibility in setting up a new partitioned index for an existing partitioned table.

- Allow foreign keys on partitioned tables (Álvaro Herrera)

- Allow FOR EACH ROW triggers on partitioned tables (Álvaro Herrera)
  Creation of a trigger on a partitioned table automatically creates triggers on all existing and future partitions. This also allows deferred unique constraints on partitioned tables.

- Allow partitioned tables to have a default partition (Jeevan Ladhe, Beena Emerson, Ashutosh Bapat, Rahila Syed, Robert Haas)
  The default partition will store rows that don't match any of the other defined partitions, and is searched accordingly.

- UPDATE statements that change a partition key column now cause affected rows to be moved to the appropriate partitions (Amit Khandekar)

- Allow INSERT, UPDATE, and COPY on partitioned tables to properly route rows to foreign partitions (Etsuro Fujita, Amit Langote)
  This is supported by postgres_fdw foreign tables.

- Allow faster partition elimination during query processing (Amit Langote, David Rowley, Dilip Kumar)
This speeds access to partitioned tables with many partitions.

- Allow partition elimination during query execution (David Rowley, Beena Emerson)

  Previously, partition elimination only happened at planning time, meaning many joins and prepared queries could not use partition elimination.

- In an equality join between partitioned tables, allow matching partitions to be joined directly (Ashutosh Bapat)

  This feature is disabled by default but can be enabled by changing enable_partition-wise_join.

- Allow aggregate functions on partitioned tables to be evaluated separately for each partition, subsequently merging the results (Jeevan Chalke, Ashutosh Bapat, Robert Haas)

  This feature is disabled by default but can be enabled by changing enable_partition-wise_aggregate.

- Allow postgres_fdw to push down aggregates to foreign tables that are partitions (Jeevan Chalke)

### E.2.3.1.2. Parallel Queries

- Allow parallel building of a btree index (Peter Geoghegan, Rushabh Lathia, Heikki Linnakangas)

- Allow hash joins to be performed in parallel using a shared hash table (Thomas Munro)

- Allow UNION to run each SELECT in parallel if the individual SELECTs cannot be parallelized (Amit Khandekar, Robert Haas, Amul Sul)

- Allow partition scans to more efficiently use parallel workers (Amit Khandekar, Robert Haas, Amul Sul)

- Allow LIMIT to be passed to parallel workers (Robert Haas, Tom Lane)

  This allows workers to reduce returned results and use targeted index scans.

- Allow single-evaluation queries, e.g. WHERE clause aggregate queries, and functions in the target list to be parallelized (Amit Kapila, Robert Haas)

- Add server parameter parallel_leader_participation to control whether the leader also executes subplans (Thomas Munro)

  The default is enabled, meaning the leader will execute subplans.

- Allow parallelization of commands CREATE TABLE ... AS, SELECT INTO, and CREATE MATERIALIZED VIEW (Haribabu Kommi)

- Improve performance of sequential scans with many parallel workers (David Rowley)

- Add reporting of parallel workers' sort activity in EXPLAIN (Robert Haas, Tom Lane)

### E.2.3.1.3. Indexes

- Allow B-tree indexes to include columns that are not part of the search key or unique constraint, but are available to be read by index-only scans (Anastasia Lubennikova, Alexander Korotkov, Teodor Sigaev)
This is enabled by the new `INCLUDE` clause of `CREATE_INDEX`. It facilitates building “covering indexes” that optimize specific types of queries. Columns can be included even if their data types don’t have B-tree support.

- Improve performance of monotonically increasing index additions (Pavan Deolasee, Peter Geoghegan)
- Improve performance of hash index scans (Ashutosh Sharma)
- Add predicate locking for hash, GiST and GIN indexes (Shubham Barai)

This reduces the likelihood of serialization conflicts in serializable-mode transactions.

### E.2.3.1.3.1. SP-Gist

- Add prefix-match operator `text ^@ text`, which is supported by SP-GiST (Ildus Kurbangaliev)

  This is similar to using `var LIKE 'word%'` with a btree index, but it is more efficient.
- Allow polygons to be indexed with SP-GiST (Nikita Glukhov, Alexander Korotkov)
- Allow SP-GiST to use lossy representation of leaf keys (Teodor Sigaev, Heikki Linnakangas, Alexander Korotkov, Nikita Glukhov)

### E.2.3.1.4. Optimizer

- Improve selection of the most common values for statistics (Jeff Janes, Dean Rasheed)

  Previously, the most common values (MCVs) were identified based on their frequency compared to all column values. Now, MCVs are chosen based on their frequency compared to the non-MCV values. This improves the robustness of the algorithm for both uniform and non-uniform distributions.
- Improve selectivity estimates for `>=` and `<=` (Tom Lane)

  Previously, such cases used the same selectivity estimates as `>` and `<`, respectively, unless the comparison constants are MCVs. This change is particularly helpful for queries involving `BETWEEN` with small ranges.
- Reduce `var = var` to `var IS NOT NULL` where equivalent (Tom Lane)

  This leads to better selectivity estimates.
- Improve optimizer’s row count estimates for `EXISTS` and `NOT EXISTS` queries (Tom Lane)
- Make the optimizer account for evaluation costs and selectivity of `HAVING` clauses (Tom Lane)

### E.2.3.1.5. General Performance

- Add Just-in-Time (JIT) compilation of some parts of query plans to improve execution speed (Andres Freund)

  This feature requires LLVM to be available. It is not currently enabled by default, even in builds that support it.
- Allow bitmap scans to perform index-only scans when possible (Alexander Kuzmenkov)
- Update the free space map during `VACUUM` (Claudio Freire)

  This allows free space to be reused more quickly.
- Allow `VACUUM` to avoid unnecessary index scans (Masahiko Sawada, Alexander Korotkov)
• Improve performance of committing multiple concurrent transactions (Amit Kapila)
• Reduce memory usage for queries using set-returning functions in their target lists (Andres Freund)
• Improve the speed of aggregate computations (Andres Freund)
• Allow postgres_fdw to push UPDATES and DELETES using joins to foreign servers (Etsuro Fujita)

Previously, only non-join UPDATES and DELETES were pushed.

• Add support for large pages on Windows (Takayuki Tsunakawa, Thomas Munro)

This is controlled by the huge_pages configuration parameter.

E.2.3.1.6. Monitoring

• Show memory usage in output from log_statement_stats, log_parser_stats, log_planner_stats, and log_executor_stats (Justin Pryzby, Peter Eisentraut)
• Add column pg_stat_activity.backend_type to show the type of a background worker (Peter Eisentraut)

The type is also visible in ps output.

• Make log_autovacuum_min_duration log skipped tables that are concurrently being dropped (Nathan Bossart)

E.2.3.1.6.1. Information Schema

• Add information_schema columns related to table constraints and triggers (Peter Eisentraut)

Specifically, triggers.action_order, triggers.action_reference_old_table, and triggers.action_reference_new_table are now populated, where before they were always null. Also, table_constraints.enforced now exists but is not yet usefully populated.

E.2.3.1.7. Authentication

• Allow the server to specify more complex LDAP specifications in search+bind mode (Thomas Munro)

Specifically, ldapssearchfilter allows pattern matching using combinations of LDAP attributes.

• Allow LDAP authentication to use encrypted LDAP (Thomas Munro)

We already supported LDAP over TLS by using ldap_tls=1. This new TLS LDAP method for encrypted LDAP is enabled with ldapscheme=ldaps or ldapurl=ldaps://.

• Improve logging of LDAP errors (Thomas Munro)

E.2.3.1.8. Permissions

• Add default roles that enable file system access (Stephen Frost)

Specifically, the new roles are: pg_read_server_files, pg_write_server_files, and pg_execute_server_program. These roles now also control who can use server-side COPY and the file_fdw extension. Previously, only superusers could use these functions, and that is still the default behavior.
- Allow access to file system functions to be controlled by GRANT/REVOKE permissions, rather than superuser checks (Stephen Frost)

Specifically, these functions were modified: pg_ls_dir(), pg_read_file(), pg_read_binary_file(), pg_stat_file().

- Use GRANT/REVOKE to control access to lo_import() and lo_export() (Michael Paquier, Tom Lane)

Previously, only superusers were granted access to these functions.

The compile-time option ALLOW_DANGEROUS_LO_FUNCTIONS has been removed.

- Use view owner not session owner when preventing non-password access to postgres_fdw tables (Robert Haas)

PostgreSQL only allows superusers to access postgres_fdw tables without passwords, e.g. via peer. Previously, the session owner had to be a superuser to allow such access; now the view owner is checked instead.

- Fix invalid locking permission check in SELECT FOR UPDATE on views (Tom Lane)

### E.2.3.1.9. Server Configuration

- Add server setting ssl_passphrase_command to allow supplying of the passphrase for SSL key files (Peter Eisentraut)

Also add ssl_passphrase_command_supports_reload to specify whether the SSL configuration should be reloaded and ssl_passphrase_command called during a server configuration reload.

- Add storage parameter toast_tuple_target to control the minimum tuple length before TOAST storage will be considered (Simon Riggs)

The default TOAST threshold has not been changed.

- Allow server options related to memory and file sizes to be specified in units of bytes (Beena Emerson)

The new unit suffix is “B”. This is in addition to the existing units “kB”, “MB”, “GB” and “TB”.

### E.2.3.1.10. Write-Ahead Log (WAL)

- Allow the WAL file size to be set during initdb (Beena Emerson)

Previously, the 16MB default could only be changed at compile time.

- Retain WAL data for only a single checkpoint (Simon Riggs)

Previously, WAL was retained for two checkpoints.

- Fill the unused portion of force-switched WAL segment files with zeros for improved compressibility (Chapman Flack)

### E.2.3.2. Base Backup and Streaming Replication

- Replicate TRUNCATE activity when using logical replication (Simon Riggs, Marco Nenciarini, Peter Eisentraut)

- Pass prepared transaction information to logical replication subscribers (Nikhil Sontakke, Stas Kelvich)
• Exclude unlogged tables, temporary tables, and pg_internal.init files from streaming base backups (David Steele)

There is no need to copy such files.

• Allow checksums of heap pages to be verified during streaming base backup (Michael Banck)

• Allow replication slots to be advanced programmatically, rather than be consumed by subscribers (Petr Jelinek)

This allows efficient advancement of replication slots when the contents do not need to be consumed. This is performed by pg_replication_slot_advance().

• Add timeline information to the backup_label file (Michael Paquier)

Also add a check that the WAL timeline matches the backup_label file's timeline.

• Add host and port connection information to the pg_stat_wal_receiver system view (Haribabu Kommi)

E.2.3.3. Utility Commands

• Allow ALTER TABLE to add a column with a non-null default without doing a table rewrite (Andrew Dunstan, Serge Rielau)

This is enabled when the default value is a constant.

• Allow views to be locked by locking the underlying tables (Yugo Nagata)

• Allow ALTER INDEX to set statistics-gathering targets for expression indexes (Alexander Korotkov, Adrien Nayrat)

In psql, \d+ now shows the statistics target for indexes.

• Allow multiple tables to be specified in one VACUUM or ANALYZE command (Nathan Bossart)

Also, if any table mentioned in VACUUM uses a column list, then the ANALYZE keyword must be supplied; previously, ANALYZE was implied in such cases.

• Add parenthesized options syntax to ANALYZE (Nathan Bossart)

This is similar to the syntax supported by VACUUM.

• Add CREATE AGGREGATE option to specify the behavior of the aggregate's finalization function (Tom Lane)

This is helpful for allowing user-defined aggregate functions to be optimized and to work as window functions.

E.2.3.4. Data Types

• Allow the creation of arrays of domains (Tom Lane)

This also allows array_agg() to be used on domains.

• Support domains over composite types (Tom Lane)

Also allow PL/Perl, PL/Python, and PL/Tcl to handle composite-domain function arguments and results. Also improve PL/Python domain handling.

• Add casts from JSONB scalars to numeric and boolean data types (Anastasia Lubennikova)
E.2.3.5. Functions

- Add all window function framing options specified by SQL:2011 (Oliver Ford, Tom Lane)

  Specifically, allow RANGE mode to use PRECEDING and FOLLOWING to select rows having grouping values within plus or minus the specified offset. Add GROUPS mode to include plus or minus the number of peer groups. Frame exclusion syntax was also added.

- Add SHA-2 family of hash functions (Peter Eisentraut)

  Specifically, sha224(), sha256(), sha384(), sha512() were added.

- Add support for 64-bit non-cryptographic hash functions (Robert Haas, Amul Sul)

- Add to_char() and to_timestamp() to specify the time zone's offset from UTC in hours and minutes (Nikita Glukhov, Andrew Dunstan)

  This is done with format specifications TZH and TZM.

- Add text search function websearch_to_tsquery() that supports a query syntax similar to that used by web search engines (Victor Drobny, Dmitry Ivanov)

- Add functions json(b)_to_tsvector() to create a text search query for matching JSON/JSONB values (Dmitry Dolgov)

E.2.3.6. Server-Side Languages

- Add SQL-level procedures, which can start and commit their own transactions (Peter Eisentraut)

  They are created with the new CREATE PROCEDURE command and invoked via CALL.

  The new ALTER/DROP ROUTINE commands allow altering/dropping of all routine-like objects, including procedures, functions, and aggregates.

  Also, writing FUNCTION is now preferred over writing PROCEDURE in CREATE OPERATOR and CREATE TRIGGER, because the referenced object must be a function not a procedure. However, the old syntax is still accepted for compatibility.

- Add transaction control to PL/pgSQL, PL/Perl, PL/Python, PL/Tcl, and SPI server-side languages (Peter Eisentraut)

  Transaction control is only available within top-transaction-level procedures and nested DO and CALL blocks that only contain other DO and CALL blocks.

- Add the ability to define PL/pgSQL composite-type variables as not null, constant, or with initial values (Tom Lane)

- Allow PL/pgSQL to handle changes to composite types (e.g. record, row) that happen between the first and later function executions in the same session (Tom Lane)

  Previously, such circumstances generated errors.

- Add extension jsonb_plpython to transform JSONB to/from PL/Python types (Anthony Bykov)

- Add extension jsonb_plperl to transform JSONB to/from PL/Perl types (Anthony Bykov)

E.2.3.7. Client Interfaces

- Change libpq to disable compression by default (Peter Eisentraut)
Compression is already disabled in modern OpenSSL versions, so that the libpq setting had no effect with such libraries.

- Add DO CONTINUE option to ecpg's WHENEVER statement (Vinayak Pokale)
  This generates a C continue statement, causing a return to the top of the contained loop when the specified condition occurs.

- Add an ecpg mode to enable Oracle Pro*C-style handling of char arrays.
  This mode is enabled with --C.

**E.2.3.8. Client Applications**

**E.2.3.8.1. psql**

- Add psql command \gdesc to display the names and types of the columns in a query result (Pavel Stehule)

- Add psql variables to report query activity and errors (Fabien Coelho)
  Specifically, the new variables are ERROR, SQLSTATE, ROW_COUNT, LAST_ERROR_MESSAGE, and LAST_ERROR_SQLSTATE.

- Allow psql to test for the existence of a variable (Fabien Coelho)
  Specifically, the syntax :(?variable_name) allows a variable's existence to be tested in an \if statement.

- Allow environment variable PSQL_PAGER to control psql's pager (Pavel Stehule)
  This allows psql's default pager to be specified as a separate environment variable from the pager for other applications. PAGER is still honored if PSQL_PAGER is not set.

- Make psql's \d+ command always show the table's partitioning information (Amit Langote, Ashutosh Bapat)
  Previously, partition information would not be displayed for a partitioned table if it had no partitions. Also indicate which partitions are themselves partitioned.

- Ensure that psql reports the proper user name when prompting for a password (Tom Lane)
  Previously, combinations of -U and a user name embedded in a URI caused incorrect reporting. Also suppress the user name before the password prompt when --password is specified.

- Allow quit and exit to exit psql when given with no prior input (Bruce Momjian)
  Also print hints about how to exit when quit and exit are used alone on a line while the input buffer is not empty. Add a similar hint for help.

- Make psql hint at using control-D when \q is entered alone on a line but ignored (Bruce Momjian)
  For example, \q does not exit when supplied in character strings.

- Improve tab completion for ALTER INDEX RESET/SET (Masahiko Sawada)

- Add infrastructure to allow psql to adapt its tab completion queries based on the server version (Tom Lane)
  Previously, tab completion queries could fail against older servers.
E.2.3.8.2. pgbench

- Add pgbench expression support for NULLs, booleans, and some functions and operators (Fabien Coelho)
- Add `\if` conditional support to pgbench (Fabien Coelho)
- Allow the use of non-ASCII characters in pgbench variable names (Fabien Coelho)
- Add pgbench option `--init-steps` to control the initialization steps performed (Masahiko Sawada)
- Add an approximately Zipfian-distributed random generator to pgbench (Alik Khilazhev)
- Allow the random seed to be set in pgbench (Fabien Coelho)
- Allow pgbench to do exponentiation with `pow()` and `power()` (Raúl Marín Rodríguez)
- Add hashing functions to pgbench (Ildar Musin)
- Make pgbench statistics more accurate when using `--latency-limit` and `--rate` (Fabien Coelho)

E.2.3.9. Server Applications

- Add an option to `pg_basebackup` that creates a named replication slot (Michael Banck)
  The option `--create-slot` creates the named replication slot (`--slot`) when the WAL streaming method (`--wal-method=stream`) is used.
- Allow initdb to set group read access to the data directory (David Steele)
  This is accomplished with the new initdb option `--allow-group-access`. Administrators can also set group permissions on the empty data directory before running initdb. Server variable `data_directory_mode` allows reading of data directory group permissions.
- Add `pg_verify_checksums` tool to verify database checksums while offline (Magnus Hagander)
- Allow `pg_resetwal` to change the WAL segment size via `--wal-segsize` (Nathan Bossart)
- Add long options to `pg_resetwal` and `pg_controldata` (Nathan Bossart, Peter Eisentraut)
- Add `pg_receivewal` option `--no-sync` to prevent synchronous WAL writes, for testing (Michael Paquier)
- Add `pg_receivewal` option `--endpos` to specify when WAL receiving should stop (Michael Paquier)
- Allow `pg_ctl` to send the SIGKILL signal to processes (Andres Freund)
  This was previously unsupported due to concerns over possible misuse.
- Reduce the number of files copied by `pg_rewind` (Michael Paquier)
- Prevent `pg_rewind` from running as root (Michael Paquier)

E.2.3.9.1. pg_dump, pg_dumpall, pg_restore

- Add `pg_dumpall` option `--encoding` to control output encoding (Michael Paquier)
• Add pg_dump option --load-via-partition-root to force loading of data into the partition's root table, rather than the original partition (Rushabh Lathia)

This is useful if the system to be loaded to has different collation definitions or endianness, possibly requiring rows to be stored in different partitions than previously.

• Add an option to suppress dumping and restoring database object comments (Robins Tharakan)

The new pg_dump, pg_dumpall, and pg_restore option is --no-comments.

E.2.3.10. Source Code

• Add PGXS support for installing include files (Andrew Gierth)

This supports creating extension modules that depend on other modules. Formerly there was no easy way for the dependent module to find the referenced one's include files. Several existing contrib modules that define data types have been adjusted to install relevant files. Also, PL/Perl and PL/Python now install their include files, to support creation of transform modules for those languages.

• Install errcodes.txt to allow extensions to access the list of error codes known to PostgreSQL (Thomas Munro)

• Convert documentation to DocBook XML (Peter Eisentraut, Alexander Lakhin, Jürgen Purtz)

The file names still use an sgml extension for compatibility with back branches.

• Use stdbool.h to define type bool on platforms where it's suitable, which is most (Peter Eisentraut)

This eliminates a coding hazard for extension modules that need to include stdbool.h.

• Overhaul the way that initial system catalog contents are defined (John Naylor)

The initial data is now represented in Perl data structures, making it much easier to manipulate mechanically.

• Prevent extensions from creating custom server parameters that take a quoted list of values (Tom Lane)

This cannot be supported at present because knowledge of the parameter's property would be required even before the extension is loaded.

• Add ability to use channel binding when using SCRAM authentication (Michael Paquier)

Channel binding is intended to prevent man-in-the-middle attacks, but SCRAM cannot prevent them unless it can be forced to be active. Unfortunately, there is no way to do that in libpq. Support for it is expected in future versions of libpq and in interfaces not built using libpq, e.g. JDBC.

• Allow background workers to attach to databases that normally disallow connections (Magnus Hagander)

• Add support for hardware CRC calculations on ARMv8 (Yuqi Gu, Heikki Linnakangas, Thomas Munro)

• Speed up lookups of built-in functions by OID (Andres Freund)

The previous binary search has been replaced by a lookup array.

• Speed up construction of query results (Andres Freund)
• Improve speed of access to system caches (Andres Freund)

• Add a generational memory allocator which is optimized for serial allocation/deallocation (Tomas Vondra)

This reduces memory usage for logical decoding.

• Make the computation of `pg_class.reltuples` by VACUUM consistent with its computation by ANALYZE (Tomas Vondra)

• Update to use perltidy version 20170521 (Tom Lane, Peter Eisentraut)

E.2.3.11. Additional Modules

• Allow extension `pg_prewarm` to restore the previous shared buffer contents on startup (Mithun Cy, Robert Haas)

This is accomplished by having `pg_prewarm` store the shared buffers' relation and block number data to disk occasionally during server operation, and at shutdown.

• Add `pg_trgm` function `strict_word_similarity()` to compute the similarity of whole words (Alexander Korotkov)

The function `word_similarity()` already existed for this purpose, but it was designed to find similar parts of words, while `strict_word_similarity()` computes the similarity to whole words.

• Allow creation of indexes that can be used by LIKE comparisons on citext columns (Alexey Chernyshov)

To do this, the index must be created using the `citext_pattern_ops` operator class.

• Allow btree_gin to index bool, bpchar, name and uuid data types (Matheus Oliveira)

• Allow `cube` and `seg` extensions to perform index-only scans using GiST indexes (Andrey Borodin)

• Allow retrieval of negative cube coordinates using the `~>` operator (Alexander Korotkov)

This is useful for KNN-GiST searches when looking for coordinates in descending order.

• Add Vietnamese letter handling to the `unaccent` extension (Dang Minh Huong, Michael Paquier)

• Enhance `amcheck` to check that each heap tuple has an index entry (Peter Geoghegan)

• Have `adminpack` use the new default file system access roles (Stephen Frost)

Previously, only superusers could call `adminpack` functions; now role permissions are checked.

• Widen `pg_stat_statement`'s query ID to 64 bits (Robert Haas)

This greatly reduces the chance of query ID hash collisions. The query ID can now potentially display as a negative value.

• Remove the contrib/start-scripts/osx scripts since they are no longer recommended (use contrib/start-scripts/macos instead) (Tom Lane)

• Remove the `chkpass` extension (Peter Eisentraut)

This extension is no longer considered to be a usable security tool or example of how to write an extension.
E.2.4. Acknowledgments

The following individuals (in alphabetical order) have contributed to this release as patch authors, committers, reviewers, testers, or reporters of issues.

Abhijit Menon-Sen
Adam Bielanski
Adam Brightwell
Adam Brusselback
Aditya Toshniwal
Adrián Escoms
Adrien Nayrat
Akos Vandra
Aleksander Alekseev
Aleksandr Parfenov
Alexander Korotkov
Alexander Kukushkin
Alexander Kuzmenkov
Alexander Lakhin
Alexandre Garcia
Alexey Bashtanov
Alexey Chernyshov
Alexey Kryuchkov
Alik Khilazhev
Álvaro Herrera
Amit Kapila
Amit Khandekar
Amit Langote
Amul Sul
Anastasia Lubennikova
Andreas Joseph Krogh
Andreas Karlsson
Andreas Seltenreich
André Hänsel
Andrei Gorita
Andres Freund
Andrew Dunstan
Andrew Fletcher
Andrew Gierth
Andrew Grossman
Andrew Krasichkov
Andrey Borodin
Andrey Lizenko
Andy Abelisto
Anthony Bykov
Antoine Scemama
Anton Dignös
Antonin Houska
Arseniy Sharoglavov
Arseny Sher
Arthur Zakirov
Ashutosh Bapat
Ashutosh Sharma
Ashwin Agrawal
Asim Praveen
Atsushi Torikoshi
Badrul Chowdhury
Balazs Szilfai
Basil Bourque
Beena Emerson
Ben Chobot
Benjamin Coutu
Bernd Helmle
Blaz Merela
Brad DeJong
Brent Dearth
Brian Cloutier
Bruce Momjian
Catalin Iacob
Chad Trabant
Chapman Flack
Christian Dutu
Christian Ullrich
Christoph Berg
Christoph Dreis
Christophe Courtois
Christopher Jones
Claudio Freire
Clayton Salem
Craig Ringer
Dagfinn Ilmari Mannsáker
Dan Vianello
Dan Watson
Dang Minh Huong
Daniel Gustafsson
Daniel Vérité
Daniel Westermann
Daniel Wood
Darafei Praliaskouski
Dave Cramer
Dave Page
David Binderman
David Carlier
David Fetter
David G. Johnston
David Gould
David Hinkle
David Pereiro Lagares
David Rader
David Rowley
David Steele
Davy Machado
Dean Rasheed
Dian Fay
Dilip Kumar
Dmitriy Sarafannikov
Dmitry Dolgov
Dmitry Ivanov
Dmitry Shalashov
Don Seiler
Doug Doole
Doug Rady
Edmund Horner
Eiji Seki
Elvis Pranskevichus
Emre Hasegeli
This release contains a variety of fixes from 10.5. For information about new features in major release 10, see Section E.9.

### E.3.1. Migration to Version 10.6

A dump/restore is not required for those running 10.X.

However, if you use the `pg_stat_statements` extension, see the changelog entry below about that.

Also, if you are upgrading from a version earlier than 10.4, see Section E.5.

### E.3.2. Changes

- Ensure proper quoting of transition table names when `pg_dump` emits `CREATE TRIGGER ... REFERENCING` commands (Tom Lane)

  This oversight could be exploited by an unprivileged user to gain superuser privileges during the next dump/reload or `pg_upgrade` run. (CVE-2018-16850)

- Fix corner-case failures in `hasFoo_privilege()` family of functions (Tom Lane)

  Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. `has_column_privilege()` was additionally capable of crashing on some platforms.
• Fix pg_get_partition_constraintdef() to return NULL rather than fail when passed an invalid relation OID (Tom Lane)

• Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)

• Fix parsing of standard multi-character operators that are immediately followed by a comment or + or − (Andrew Gierth)

This oversight could lead to parse errors, or to incorrect assignment of precedence.

• Avoid O(N^3) slowdown in lexer for long strings of + or − characters (Andrew Gierth)

• Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)

• Fix failure of UPDATE/DELETE ... WHERE CURRENT OF ... after rewinding the referenced cursor (Tom Lane)

A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

• Fix EvalPlanQual to handle conditionally-executed InitPlans properly (Andrew Gierth, Tom Lane)

This resulted in hard-to-reproduce crashes or wrong answers in concurrent updates, if they contained code such as an uncorrelated sub-SELECT inside a CASE construct.

• Prevent creation of a partition in a trigger attached to its parent table (Amit Langote)

Ideally we’d allow that, but for the moment it has to be blocked to avoid crashes.

• Fix problems with applying ON COMMIT DELETE ROWS to a partitioned temporary table (Amit Langote)

• Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)

This bug affected full-text-search operations, as well as contrib/ltree and contrib/pg_trgm.

• Disallow pushing sub-SELECTs containing window functions, LIMIT, or OFFSET to parallel workers (Amit Kapila)

Such cases could result in inconsistent behavior due to different workers getting different answers, as a result of indeterminacy due to row-ordering variations.

• Ensure that sequences owned by a foreign table are processed by ALTER OWNER on the table (Peter Eisentraut)

The ownership change should propagate to such sequences as well, but this was missed for foreign tables.

• Ensure that the server will process already-received NOTIFY and SIGTERM interrupts before waiting for client input (Jeff Janes, Tom Lane)

• Fix over-allocation of space for array_out()’s result string (Keiichi Hirobe)

• Avoid query-lifetime memory leak in XMLTABLE (Andrew Gierth)

• Fix memory leak in repeated SP-GiST index scans (Tom Lane)
This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- Ensure that `ApplyLogicalMappingFile()` closes the mapping file when done with it (Tomas Vondra)

  Previously, the file descriptor was leaked, eventually resulting in failures during logical decoding.

- Fix logical decoding to handle cases where a mapped catalog table is repeatedly rewritten, e.g. by `VACUUM FULL` (Andres Freund)

- Prevent starting the server with `wal_level` set to too low a value to support an existing replication slot (Andres Freund)

- Avoid crash if a utility command causes infinite recursion (Tom Lane)

- When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master (Michael Paquier, Konstantin Knizhnik)

- Fix event triggers to handle nested `ALTER TABLE` commands (Michael Paquier, Álvaro Herrera)

- Propagate parent process's transaction and statement start timestamps to parallel workers (Konstantin Knizhnik)

  This prevents misbehavior of functions such as `transaction_timestamp()` when executed in a worker.

- Fix transfer of expanded datums to parallel workers so that alignment is preserved, preventing crashes on alignment-picky platforms (Tom Lane, Amit Kapila)

- Fix WAL file recycling logic to work correctly on standby servers (Michael Paquier)

  Depending on the setting of `archive_mode`, a standby might fail to remove some WAL files that could be removed.

- Fix handling of commit-timestamp tracking during recovery (Masahiko Sawasa, Michael Paquier)

  If commit timestamp tracking has been turned on or off, recovery might fail due to trying to fetch the commit timestamp for a transaction that did not record it.

- Randomize the `random()` seed in bootstrap and standalone backends, and in `initdb` (Noah Misch)

  The main practical effect of this change is that it avoids a scenario where `initdb` might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

- Fix possible shared-memory corruption in DSA logic (Thomas Munro)

- Allow DSM allocation to be interrupted (Chris Travers)

- Avoid failure in a parallel worker when loading an extension that tries to access system caches within its init function (Thomas Munro)

  We don't consider that to be good extension coding practice, but it mostly worked before parallel query, so continue to support it for now.

- Properly handle turning `full_page_writes` on dynamically (Kyotaro Horiguchi)

- Fix possible crash due to double `free()` during SP-GiST rescan (Andrew Gierth)

- Prevent mis-linking of `src/port` and `src/common` functions on ELF-based BSD platforms, as well as HP-UX and Solaris (Andrew Gierth, Tom Lane)
Shared libraries loaded into a backend’s address space could use the backend’s versions of these functions, rather than their own copies as intended. Since the behavior of the two sets of functions isn’t quite the same, this led to failures.

- Avoid possible buffer overrun when replaying GIN page recompression from WAL (Alexander Korotkov, Sivasubramanian Ramasubramanian)
- Avoid overrun of a hash index’s metapage when BLCKSZ is smaller than default (Dilip Kumar)
- Fix missed page checksum updates in hash indexes (Amit Kapila)
- Fix missed fsync of a replication slot’s directory (Konstantin Knizhnik, Michael Paquier)
- Fix unexpected timeouts when using wal_sender_timeout on a slow server (Noah Misch)
- Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)

This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.

- Ensure background workers are stopped properly when the postmaster receives a fast-shutdown request before completing database startup (Alexander Kukushkin)
- Update the free space map during WAL replay of page all-visible/frozen flag changes (Álvaro Herrera)

Previously we were not careful about this, reasoning that the FSM is not critical data anyway. However, if it’s sufficiently out of date, that can result in significant performance degradation after a standby has been promoted to primary. The FSM will eventually be healed by updates, but we’d like it to be good sooner, so work harder at maintaining it during WAL replay.

- Avoid premature release of parallel-query resources when query end or tuple count limit is reached (Amit Kapila)

It’s only okay to shut down the executor at this point if the caller cannot demand backwards scan afterwards.

- Don’t run atexit callbacks when servicing SIGQUIT (Heikki Linnakangas)
- Don’t record foreign-server user mappings as members of extensions (Tom Lane)

If CREATE USER MAPPING is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can’t be extension members, so user mappings shouldn’t be either.

- Make syslogger more robust against failures in opening CSV log files (Tom Lane)
- When libpq is given multiple target host names, do the DNS lookups one at a time, not all at once (Tom Lane)

This prevents unnecessary failures or slow connections when a connection is successfully made to one of the earlier servers in the list.

- Fix libpq’s handling of connection timeouts so that they are properly applied per host name or IP address (Tom Lane)

Previously, some code paths failed to restart the timer when switching to a new target host, possibly resulting in premature timeout.
• Fix psql, as well as documentation examples, to call PQconsumeInput() before each NOTIFY call (Tom Lane)

This fixes cases in which psql would not report receipt of a NOTIFY message until after the next command.

• Fix pg_dump's --no-publications option to also ignore publication tables (Gilles Darold)

• In pg_dump, exclude identity sequences when their parent table is excluded from the dump (David Rowley)

• Fix possible inconsistency in pg_dump's sorting of dissimilar object names (Jacob Champion)

• Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)

This avoids failures due to the new policy of running restores with restrictive search path.

• Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)

pg_upgrade failed to preserve an event trigger's extension-membership status.

• Fix pg_upgrade's cluster state check to work correctly on a standby server (Bruce Momjian)

• Enforce type cube's dimension limit in all contrib/cube functions (Andrey Borodin)

Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.

• In contrib/pg_stat_statements, disallow the pg_read_all_stats role from executing pg_stat_statements_reset() (Haribabu Kommi)

pg_read_all_stats is only meant to grant permission to read statistics, not to change them, so this grant was incorrect.

To cause this change to take effect, run ALTER EXTENSION pg_stat_statements UPDATE in each database where pg_stat_statements has been installed.

• In contrib/postgres_fdw, don't try to ship a variable-free ORDER BY clause to the remote server (Andrew Gierth)

• Fix contrib/unaccent's unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)

Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.

• Fix build problems on macOS 10.14 (Mojave) (Tom Lane)

Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.

It is now recommended that Perl-related extensions write $(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.

Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.

• Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)
Perl no longer includes the current directory in its search path by default; work around that.

- On Windows, allow the regression tests to be run by an Administrator account (Andrew Dunstan)
  To do this safely, pg_regress now gives up any such privileges at startup.
- Allow btree comparison functions to return INT_MIN (Tom Lane)
  Up to now, we've forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of memcmp(), strcmp(), etc, as POSIX doesn't place any such restriction on those functions. At least some recent versions ofmemcmp() can return INT_MIN, causing incorrect sort ordering. Hence, we've removed this restriction. Callers must now use the INVERT_COMPARE_RESULT() macro if they wish to invert the sort order.
- Fix recursion hazard in shared-invalidation message processing (Tom Lane)
  This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.
  This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.
- Save and restore SPI's global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)
  This prevents possible interference when one SPI-using function calls another.
- Avoid using potentially-under-aligned page buffers (Tom Lane)
  Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can't place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.
- Make src/port/snprintf.c follow the C99 standard's definition of snprintf()'s result value (Tom Lane)
  On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.
- When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)
  This avoids problems with missed floating point overflow checks.
- Fix configure's detection of the result type of strerror_r() (Tom Lane)
  The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.
- Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.4. Release 10.5

Release date: 2018-08-09

This release contains a variety of fixes from 10.4. For information about new features in major release 10, see Section E.9.
E.4.1. Migration to Version 10.5

A dump/restore is not required for those running 10.X.

However, if you are upgrading from a version earlier than 10.4, see Section E.5.

E.4.2. Changes

• Fix failure to reset libpq’s state fully between connection attempts (Tom Lane)

An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general, an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasichkov for reporting this issue. (CVE-2018-10915)

• Fix INSERT ... ON CONFLICT UPDATE through a view that isn’t just SELECT * FROM ...
  (Dean Rasheed, Amit Langote)

Erroneous expansion of an updatable view could lead to crashes or “attribute ... has the wrong type” errors, if the view’s SELECT list doesn’t match one-to-one with the underlying table’s columns. Furthermore, this bug could be leveraged to allow updates of columns that an attacking user lacks UPDATE privilege for, if that user has INSERT and UPDATE privileges for some other column(s) of the table. Any user could also use it for disclosure of server memory. (CVE-2018-10925)

• Ensure that updates to the relfrozenxid and relminmxid values for “nailed” system catalogs are processed in a timely fashion (Andres Freund)

Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as pg_authid, because the stale cache data could persist into new sessions as well as existing ones.

• Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

• Avoid emitting a bogus WAL record when recycling an all-zero btree page (Amit Kapila)

This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

• During WAL replay, guard against corrupted record lengths exceeding 1GB (Michael Paquier)

Treat such a case as corrupt data. Previously, the code would try to allocate space and get a hard error, making recovery impossible.

• When ending recovery, delay writing the timeline history file as long as possible (Heikki Linnakan-gas)

This avoids some situations where a failure during recovery cleanup (such as a problem with a two-phase state file) led to inconsistent timeline state on-disk.

• Improve performance of WAL replay for transactions that drop many relations (Fujii Masao)
This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

- Improve performance of lock releasing in standby server WAL replay (Thomas Munro)
- Make logical WAL senders report streaming state correctly (Simon Riggs, Sawada Masahiko)
  The code previously mis-detected whether or not it had caught up with the upstream server.
- Ensure that a snapshot is provided when executing data type input functions in logical replication subscribers (Minh-Quan Tran, Álvaro Herrera)
  This omission led to failures in some cases, such as domains with constraints using SQL-language functions.
- Fix bugs in snapshot handling during logical decoding, allowing wrong decoding results in rare cases (Arseny Sher, Álvaro Herrera)
- Add subtransaction handling in logical-replication table synchronization workers (Amit Khandekar, Robert Haas)
  Previously, table synchronization could misbehave if any subtransactions were aborted after modifying a table being synchronized.
- Ensure a table’s cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)
  Previously, the failed index’s OID could remain in the list, causing problems later in the same session.
- Fix mishandling of empty uncompressed posting list pages in GIN indexes (Sivasubramanian Ramasubramanian, Alexander Korotkov)
  This could result in an assertion failure after pg_upgrade of a pre-9.4 GIN index (9.4 and later will not create such pages).
- Pad arrays of unnamed POSIX semaphores to reduce cache line sharing (Thomas Munro)
  This reduces contention on many-CPU systems, fixing a performance regression (compared to previous releases) on Linux and FreeBSD.
- Ensure that a process doing a parallel index scan will respond to signals (Amit Kapila)
  Previously, parallel workers could get stuck waiting for a lock on an index page, and not notice requests to abort the query.
- Ensure that VACUUM will respond to signals within btree page deletion loops (Andres Freund)
  Corrupted btree indexes could result in an infinite loop here, and that previously wasn’t interruptible without forcing a crash.
- Fix hash-join costing mistake introduced with inner_unique optimization (David Rowley)
  This could lead to bad plan choices in situations where that optimization was applicable.
- Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)
  This resulted in failure to recognize that an index on a composite column could provide the sort order needed for a mergejoin on that column.
- Fix planner to avoid “ORDER/GROUP BY expression not found in targetlist” errors in some queries with set-returning functions (Tom Lane)
• Fix handling of partition keys whose data type uses a polymorphic btree operator class, such as arrays (Amit Langote, Álvaro Herrera)

• Fix SQL-standard `FETCH FIRST` syntax to allow parameters (Andrew Gierth)

• Remove undocumented restriction against duplicate partition key columns (Yugo Nagata)

• Disallow temporary tables from being partitions of non-temporary tables (Amit Langote, Michael Paquier)

While previously allowed, this case didn’t work reliably.

• Fix `EXPLAIN`’s accounting for resource usage, particularly buffer accesses, in parallel workers (Amit Kapila, Robert Haas)

• Fix `SHOW ALL` to show all settings to roles that are members of `pg_read_all_settings`, and also allow such roles to see source filename and line number in the `pg_settings` view (Laurenz Albe, Álvaro Herrera)

• Fix failure to schema-qualify some object names in `getObjectDescription` and `getObjectIdentity` output (Kyotaro Horiguchi, Tom Lane)

Names of collations, conversions, text search objects, publication relations, and extended statistics objects were not schema-qualified when they should be.

• Fix `CREATE AGGREGATE` type checking so that parallelism support functions can be attached to variadic aggregates (Alexey Bashtanov)

• Widen `COPY FROM`’s current-line-number counter from 32 to 64 bits (David Rowley)

This avoids two problems with input exceeding 4G lines: `COPY FROM WITH HEADER` would drop a line every 4G lines, not only the first line, and error reports could show a wrong line number.

• Allow replication slots to be dropped in single-user mode (Álvaro Herrera)

This use-case was accidentally broken in release 10.0.

• Fix incorrect results from `variance(int4)` and related aggregates when run in parallel aggregation mode (David Rowley)

• Process `TEXT` and `CDATA` nodes correctly in `xmltable()` column expressions (Markus Winand)

• Cope with possible failure of OpenSSL’s `RAND_bytes()` function (Dean Rasheed, Michael Paquier)

Under rare circumstances, this oversight could result in “could not generate random cancel key” failures that could only be resolved by restarting the postmaster.

• Fix `libpq`’s handling of some cases where `hostaddr` is specified (Hari Babu, Tom Lane, Robert Haas)

`PQhost()` gave misleading or incorrect results in some cases. Now, it uniformly returns the host name if specified, or the host address if only that is specified, or the default host name (typically `/tmp` or `localhost`) if both parameters are omitted.

Also, the wrong value might be compared to the server name when verifying an SSL certificate.

Also, the wrong value might be compared to the host name field in `~/.pgpass`. Now, that field is compared to the host name if specified, or the host address if only that is specified, or `localhost` if both parameters are omitted.
Also, an incorrect error message was reported for an unparsable hostaddr value.

Also, when the host, hostaddr, or port parameters contain comma-separated lists, libpq is now more careful to treat empty elements of a list as selecting the default behavior.

• Add a string freeing function to ecpg’s pgtypes library, so that cross-module memory management problems can be avoided on Windows (Takayuki Tsunakawa)

On Windows, crashes can ensue if the free call for a given chunk of memory is not made from the same DLL that malloc’d the memory. The pgtypes library sometimes returns strings that it expects the caller to free, making it impossible to follow this rule. Add a PGTYPESchar_free() function that just wraps free, allowing applications to follow this rule.

• Fix ecpg's support for long long variables on Windows, as well as other platforms that declare strtoll/strtoull nonstandardly or not at all (Dang Minh Huong, Tom Lane)

• Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in the semantics of a statement intra-session (Tom Lane)

This error led to assertion failures, or in rare cases, failure to enforce the INTO STRICT option as expected.

• Fix password prompting in client programs so that echo is properly disabled on Windows when stdin is not the terminal (Matthew Stickney)

• Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)

The previous fix for quoting of search_path and other list-valued variables in pg_dump output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

• Fix pg_dump's failure to dump REPLICA IDENTITY properties for constraint indexes (Tom Lane)

Manually created unique indexes were properly marked, but not those created by declaring UNIQUE or PRIMARY KEY constraints.

• Make pg_upgrade check that the old server was shut down cleanly (Bruce Momjian)

The previous check could be fooled by an immediate-mode shutdown.

• Fix contrib/hstore_plperl to look through Perl scalar references, and to not crash if it doesn’t find a hash reference where it expects one (Tom Lane)

• Fix crash in contrib/ltree’s lca() function when the input array is empty (Pierre Ducroquet)

• Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

• Rearrange makefiles to ensure that programs link to freshly-built libraries (such as libpq.so) rather than ones that might exist in the system library directories (Tom Lane)

This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

• Update time zone data files to tzdata release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.

This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbreviations in use; the only known effect is that the is_dst column in the pg_timezone_names view will now be true in winter and false in summer in these cases.
**E.5. Release 10.4**

**Release date:** 2018-05-10

This release contains a variety of fixes from 10.3. For information about new features in major release 10, see Section E.9.

**E.5.1. Migration to Version 10.4**

A dump/restore is not required for those running 10.X.

However, if you use the `adminpack` extension, you should update it as per the first changelog entry below.

Also, if the function marking mistakes mentioned in the second and third changelog entries below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 10.3, see Section E.6.

**E.5.2. Changes**

- Remove public execute privilege from `contrib/adminpack`'s `pg_logfile_rotate()` function (Stephen Frost)

  `pg_logfile_rotate()` is a deprecated wrapper for the core function `pg_rotate_logfile()`. When that function was changed to rely on SQL privileges for access control rather than a hard-coded superuser check, `pg_logfile_rotate()` should have been updated as well, but the need for this was missed. Hence, if `adminpack` is installed, any user could request a logfile rotation, creating a minor security issue.

  After installing this update, administrators should update `adminpack` by performing `ALTER EXTENSION adminpack UPDATE` in each database in which `adminpack` is installed.

  (CVE-2018-1115)

- Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

  The functions `query_to_xml`, `cursor_to_xml`, `cursor_to_xmlschema`, `query_to_xmlschema`, and `query_to_xml_and_xmlschema` should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions' `pg_proc` entries, for example `ALTER FUNCTION pg_catalog.query_to_xml(text, boolean, boolean, text) VOLATILE`.

- Fix incorrect parallel-safety markings on a few built-in functions (Thomas Munro, Tom Lane)

  The functions `brin_summarize_new_values`, `brin_summarize_range`, `brin_desummarize_range`, `gin_clean_pending_list`, `cursor_to_xml`, `cursor_to_xmlschema`, `ts_rewrite`, `ts_stat`, `binary_upgrade_create_empty_extension`, and `pg_import_system_collations` should be marked parallel-unsafe; some because they perform database modifications directly, and others because they execute user-supplied queries that might do so. They were marked parallel-restricted instead, leading to a risk of unexpected query errors. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard unless `force_parallel_mode` is turned on. In
case of trouble, it can be fixed by manually updating these functions' \texttt{pg\_procedures} entries, for example
\texttt{ALTER FUNCTION pg\_catalog.brin\_summarize\_new\_values(regclass) PAR\-ALLEL UNSAFE}. (Note that that will need to be done in each database of the installation.) Another option is to \texttt{pg\_upgrade} the database to a version containing the corrected initial data.

- Avoid re-using TOAST value OIDs that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

Once the OID counter has wrapped around, it's possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value \texttt{nunnun}” errors, which would persist until the dead entry was removed by \texttt{VACUUM}. Fix by not selecting such OIDs when creating a new TOAST entry.

- Correctly enforce any \texttt{CHECK} constraints on individual partitions during \texttt{COPY} to a partitioned table (Etsuro Fujita)

Previously, only constraints declared for the partitioned table as a whole were checked.

- Accept \texttt{TRUE} and \texttt{FALSE} as partition bound values (Amit Langote)

Previously, only string-literal values were accepted for a boolean partitioning column. But then \texttt{pg\_dump} would print such values as \texttt{TRUE} or \texttt{FALSE}, leading to dump/reload failures.

- Fix memory management for partition key comparison functions (Álvaro Herrera, Amit Langote)

This error could lead to crashes when using user-defined operator classes for partition keys.

- Fix possible crash when a query inserts tuples in several partitions of a partitioned table, and those partitions don't have identical row types (Etsuro Fujita, Amit Langote)

- Change \texttt{ANALYZE}'s algorithm for updating \texttt{pg\_class}\_reltuples (David Gould)

Previously, pages not actually scanned by \texttt{ANALYZE} were assumed to retain their old tuple density. In a large table where \texttt{ANALYZE} samples only a small fraction of the pages, this meant that the overall tuple density estimate could not change very much, so that \texttt{reltuples} would change nearly proportionally to changes in the table's physical size (\texttt{relpages}) regardless of what was actually happening in the table. This has been observed to result in \texttt{reltuples} becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that \texttt{ANALYZE}'s sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

- Include extended-statistics objects in the set of table properties duplicated by \texttt{CREATE TABLE \ldots LIKE \ldots INCLUDING ALL} (David Rowley)

Also add an \texttt{INCLUDING STATISTICS} option, to allow finer-grained control over whether this happens.

- Fix \texttt{CREATE TABLE \ldots LIKE} with \texttt{bigint} identity columns (Peter Eisentraut)

On platforms where \texttt{long} is 32 bits (which includes 64-bit Windows as well as most 32-bit machines), copied sequence parameters would be truncated to 32 bits.

- Avoid deadlocks in concurrent \texttt{CREATE INDEX CONCURRENTLY} commands that are run under \texttt{SERIALIZABLE} or \texttt{REPEATABLE READ} transaction isolation (Tom Lane)

- Fix possible slow execution of \texttt{REFRESH MATERIALIZED VIEW CONCURRENTLY} (Thomas Munro)

- Fix \texttt{UPDATE/DELETE \ldots WHERE CURRENT} to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)
• Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)
  This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.

• Fix possibly incorrect generation of an index-only-scan plan when the same table column appears in multiple index columns, and only some of those index columns use operator classes that can return the column value (Kyotaro Horiguchi)

• Fix misoptimization of CHECK constraints having provably-NULL subclauses of top-level AND/OR conditions (Tom Lane, Dean Rasheed)
  This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.

• Prevent planner crash when a query has multiple GROUPING SETS, none of which can be implemented by sorting (Andrew Gierth)

• Fix executor crash due to double free in some GROUPING SETS usages (Peter Geoghegan)

• Fix misexecution of self-joins on transition tables (Thomas Munro)

• Avoid crash if a table rewrite event trigger is added concurrently with a command that could call such a trigger (Álvaro Herrera, Andrew Gierth, Tom Lane)

• Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)

• Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)

• Fix possible leak or double free of visibility map buffer pins (Amit Kapila)

• Avoid spuriously marking pages as all-visible (Dan Wood, Pavan Deolasee, Álvaro Herrera)
  This could happen if some tuples were locked (but not deleted). While queries would still function correctly, vacuum would normally ignore such pages, with the long-term effect that the tuples were never frozen. In recent releases this would eventually result in errors such as “found multixact nnnnn from before relminmxid nnnnn”.

• Fix overly strict sanity check in heap_prepare_freeze_tuple (Álvaro Herrera)
  This could result in incorrect “cannot freeze committed xmax” failures in databases that have been pg_upgrade'd from 9.2 or earlier.

• Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)

• Reduce locking during autovacuum worker scheduling (Jeff Janes)
  The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.

• Ensure client hostname is copied while copying pg_stat_activity data to local memory (Edmund Horner)
  Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the client hostname column to change unexpectedly if any existing session disconnected.

• Handle pg_stat_activity information for auxiliary processes correctly (Edmund Horner)
  The application_name, client_hostname, and query fields might show incorrect data for such processes.
Release Notes

- Fix incorrect processing of multiple compound affixes in `ispell` dictionaries (Arthur Zakirov)
- Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on text columns (Tom Lane)

  Such searches would return the wrong set of rows in most non-C locales.
- Prevent query-lifespan memory leakage with SP-GiST operator classes that use traversal values (Anton Dignös)
- Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)

  Previously, the tuple count was reported to be the same as that of the underlying table, which is wrong if the index is partial.
- Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)

  Previously it reported the estimated number of heap tuples, which might be inaccurate, and is certainly wrong if the index is partial.
- Fix a corner case where a streaming standby gets stuck at a WAL continuation record (Kyotaro Horiguchi)
- In logical decoding, avoid possible double processing of WAL data when a walsender restarts (Craig Ringer)
- Fix logical replication to not assume that type OIDs match between the local and remote servers (Masahiko Sawada)
- Allow `scalarltsel` and `scalargtsel` to be used on non-core datatypes (Tomas Vondra)
- Reduce libpq's memory consumption when a server error is reported after a large amount of query output has been collected (Tom Lane)

  Discard the previous output before, not after, processing the error message. On some platforms, notably Linux, this can make a difference in the application's subsequent memory footprint.
- Fix double-free crashes in `ecpg` (Patrick Krecker, Jeevan Ladhe)
- Fix `ecpg` to handle `long long int` variables correctly in MSVC builds (Michael Meskes, Andrew Gierth)
- Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)

  The `local_preload_libraries`, `session_preload_libraries`, `shared_preload_libraries`, and `temp_tablespaces` variables were not correctly quoted in `pg_dump` output. This would cause problems if settings for these variables appeared in `CREATE FUNCTION ... SET` or `ALTER DATABASE/ROLE ... SET` clauses.
- Fix `pg_recvlogical` to not fail against pre-v10 PostgreSQL servers (Michael Paquier)

  A previous fix caused `pg_recvlogical` to issue a command regardless of server version, but it should only be issued to v10 and later servers.
- Ensure that `pg_rewind` deletes files on the target server if they are deleted from the source server during the run (Takayuki Tsunakawa)

  Failure to do this could result in data inconsistency on the target, particularly if the file in question is a WAL segment.
- Fix `pg_rewind` to handle tables in non-default tablespaces correctly (Takayuki Tsunakawa)
• Fix overflow handling in PL/pgSQL integer FOR loops (Tom Lane)

The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

• Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisentraut)

• Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)

• Fix errors in initial build of contrib/bloom indexes (Tomas Vondra, Tom Lane)

Fix possible omission of the table's last tuple from the index. Count the number of index tuples correctly, in case it is a partial index.

• Rename internal b64_encode and b64_decode functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)

• Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)

This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.

• Update time zone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

E.6. Release 10.3

Release date: 2018-03-01

This release contains a variety of fixes from 10.2. For information about new features in major release 10, see Section E.9.

E.6.1. Migration to Version 10.3

A dump/restore is not required for those running 10.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that you read the documentation changes described in the first changelog entry below, and take suitable steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 10.2, see Section E.7.

E.6.2. Changes

• Document how to configure installations and applications to guard against search-path-dependent trojan-horse attacks from other users (Noah Misch)

Using a search_path setting that includes any schemas writable by a hostile user enables that user to capture control of queries and then run arbitrary SQL code with the permissions of the attacked user. While it is possible to write queries that are proof against such hijacking, it is notation-
ally tedious, and it's very easy to overlook holes. Therefore, we now recommend configurations in which no untrusted schemas appear in one's search path. Relevant documentation appears in Section 5.8.6 (for database administrators and users), Section 34.1 (for application authors), Section 38.16.1 (for extension authors), and CREATE FUNCTION (for authors of SECURITY DEFINER functions). (CVE-2018-1058)

- Avoid use of insecure search_path settings in pg_dump and other client programs (Noah Misch, Tom Lane)

pg_dump, pg_upgrade, vacuumdb and other PostgreSQL-provided applications were themselves vulnerable to the type of hijacking described in the previous changelog entry; since these applications are commonly run by superusers, they present particularly attractive targets. To make them secure whether or not the installation as a whole has been secured, modify them to include only the pg_catalog schema in their search_path settings. Autovacuum worker processes now do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example, user-provided functions in index expressions — the tighter search_path may result in errors, which will need to be corrected by adjusting those user-provided functions to not assume anything about what search path they are invoked under. That has always been good practice, but now it will be necessary for correct behavior. (CVE-2018-1058)

- Prevent logical replication from trying to ship changes for unpublishable relations (Peter Eisentraut)

A publication marked FOR ALL TABLES would incorrectly ship changes in materialized views and information_schema tables, which are supposed to be omitted from the change stream.

- Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom Lane)

If a CTE (WITH clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

- Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do not match mergeclauses” planner errors in corner cases.

- Repair pg_upgrade’s failure to preserve relfrozenxid for materialized views (Tom Lane, Andres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The problem would be more likely to occur in seldom-refreshed materialized views, or ones that were maintained only with REFRESH MATERIALIZED VIEW CONCURRENTLY.

If such corruption is observed, it can be repaired by refreshing the materialized view (without CONCURRENTLY).

- Fix incorrect pg_dump output for some non-default sequence limit values (Alexey Bashtanov)

- Fix pg_dump's mishandling of STATISTICS objects (Tom Lane)

An extended statistics object's schema was mislabeled in the dump's table of contents, possibly leading to the wrong results in a schema-selective restore. Its ownership was not correctly restored, either. Also, change the logic so that statistics objects are dumped/restored, or not, as independent objects rather than tying them to the dump/restore decision for the table they are on. The original definition could not scale to the planned future extension to cross-table statistics.

- Fix incorrect reporting of PL/Python function names in error CONTEXT stacks (Tom Lane)
An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function's name twice, rather than the expected results. Also, an error in a nested PL/Python `DO` block could result in a null pointer dereference crash on some platforms.

- Allow `contrib/auto_explain's log_min_duration` setting to range up to `INT_MAX`, or about 24 days instead of 35 minutes (Tom Lane)

- Mark assorted GUC variables as `PGDLLIMPORT`, to ease porting extension modules to Windows (Metin Doslu)

## E.7. Release 10.2

**Release date:** 2018-02-08

This release contains a variety of fixes from 10.1. For information about new features in major release 10, see Section E.9.

### E.7.1. Migration to Version 10.2

A dump/restore is not required for those running 10.X.

However, if you use `contrib/cube`'s `~>` operator, see the entry below about that.

Also, if you are upgrading from a version earlier than 10.1, see Section E.8.

### E.7.2. Changes

- Fix processing of partition keys containing multiple expressions (Álvaro Herrera, David Rowley)

  This error led to crashes or, with carefully crafted input, disclosure of arbitrary backend memory. (CVE-2018-1052)

- Ensure that all temporary files made by `pg_upgrade` are non-world-readable (Tom Lane, Noah Misch)

  `pg_upgrade` normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing `pg_dumpall -g` output would be group- or world-readable, or even writable, if the user's `umask` setting allows. In typical usage on multi-user machines, the `umask` and/or the working directory's permissions would be tight enough to prevent problems; but there may be people using `pg_upgrade` in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

- Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

  In some cases `VACUUM` would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

- Fix failure to mark a hash index's metapage dirty after adding a new overflow page, potentially leading to index corruption (Lixian Zou, Amit Kapila)

- Ensure that vacuum will always clean up the pending-insertions list of a GIN index (Masahiko Sawada)

  This is necessary to ensure that dead index entries get removed. The old code got it backwards, allowing vacuum to skip the cleanup if some other process were running cleanup concurrently, thus risking invalid entries being left behind in the index.
- Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)
  These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.
- Fix incorrect query results from cases involving flattening of subqueries whose outputs are used in GROUPING SETS (Heikki Linnakangas)
- Fix handling of list partitioning constraints for partition keys of boolean or array types (Amit Langote)
- Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by ALTER TABLE NO INHERIT (Tom Lane)
- Fix spurious deadlock failures when multiple sessions are running CREATE INDEX CONCURRENTLY (Jeff Janes)
- During VACUUM FULL, update the table's size fields in pg_class sooner (Amit Kapila)
  This prevents poor behavior when rebuilding hash indexes on the table, since those use the pg_class statistics to govern the initial hash size.
- Fix UNION/INTERSECT/EXCEPT over zero columns (Tom Lane)
- Disallow identity columns on typed tables and partitions (Michael Paquier)
  These cases will be treated as unsupported features for now.
- Fix assorted failures to apply the correct default value when inserting into an identity column (Michael Paquier, Peter Eisentraut)
  In several contexts, notably COPY and ALTER TABLE ADD COLUMN, the expected default value was not applied and instead a null value was inserted.
- Fix failures when an inheritance tree contains foreign child tables (Etsuro Fujita)
  A mix of regular and foreign tables in an inheritance tree resulted in creation of incorrect plans for UPDATE and DELETE queries. This led to visible failures in some cases, notably when there are row-level triggers on a foreign child table.
- Repair failure with correlated sub-SELECT inside VALUES inside a LATERAL subquery (Tom Lane)
- Fix "could not devise a query plan for the given query" planner failure for some cases involving nested UNION ALL inside a lateral subquery (Tom Lane)
- Allow functional dependency statistics to be used for boolean columns (Tom Lane)
  Previously, although extended statistics could be declared and collected on boolean columns, the planner failed to apply them.
- Avoid underestimating the number of groups emitted by subqueries containing set-returning functions in their grouping columns (Tom Lane)
  Cases similar to SELECT DISTINCT unnest(foo) got a lower output rowcount estimate in 10.0 than they did in earlier releases, possibly resulting in unfavorable plan choices. Restore the prior estimation behavior.
- Fix use of triggers in logical replication workers (Petr Jelinek)
- Fix logical decoding to correctly clean up disk files for crashed transactions (Atsushi Torikoshi)
Logical decoding may spill WAL records to disk for transactions generating many WAL records. Normally these files are cleaned up after the transaction's commit or abort record arrives; but if no such record is ever seen, the removal code misbehaved.

- Fix walsender timeout failure and failure to respond to interrupts when processing a large transaction (Petr Jelinek)
- Fix race condition during replication origin drop that could allow the dropping process to wait indefinitely (Tom Lane)
- Allow members of the pg_read_all_stats role to see walsender statistics in the pg_stat_replication view (Feike Steenbergen)
- Show walsenders that are sending base backups as active in the pg_stat_activity view (Magnus Hagander)
- Fix reporting of scram-sha-256 authentication method in the pg_hba_file_rules view (Michael Paquier)

Previously this was printed as scram-sha256, possibly confusing users as to the correct spelling.

- Fix has_sequence_privilege() to support WITH GRANT OPTION tests, as other privilege-testing functions do (Joe Conway)

- In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)

We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don't promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only xpath() and related functions; other XML code paths already acted this way.

- Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)

Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.

- Allow a client that supports SCRAM channel binding (such as v11 or later libpq) to connect to a v10 server (Michael Paquier)

v10 does not have this feature, and the connection-time negotiation about whether to use it was done incorrectly.

- Avoid live-lock in ConditionVariableBroadcast() (Tom Lane, Thomas Munro)

Given repeatedly-unlucky timing, a process attempting to awaken all waiters for a condition variable could loop indefinitely. Due to the limited usage of condition variables in v10, this affects only parallel index scans and some operations on replication slots.

- Clean up waits for condition variables correctly during subtransaction abort (Robert Haas)

- Ensure that child processes that are waiting for a condition variable will exit promptly if the postmaster process dies (Tom Lane)

- Fix crashes in parallel queries using more than one Gather node (Thomas Munro)
• Fix hang in parallel index scan when processing a deleted or half-dead index page (Amit Kapila)

• Avoid crash if parallel bitmap heap scan is unable to allocate a shared memory segment (Robert Haas)

• Cope with failure to start a parallel worker process (Amit Kapila, Robert Haas)

  Parallel query previously tended to hang indefinitely if a worker could not be started, as the result of fork() failure or other low-probability problems.

• Avoid unnecessary failure when no parallel workers can be obtained during parallel query startup (Robert Haas)

• Fix collection of EXPLAIN statistics from parallel workers (Amit Kapila, Thomas Munro)

• Ensure that query strings passed to parallel workers are correctly null-terminated (Thomas Munro)

  This prevents emitting garbage in postmaster log output from such workers.

• Avoid unsafe alignment assumptions when working with __int128 (Tom Lane)

  Typically, compilers assume that __int128 variables are aligned on 16-byte boundaries, but our memory allocation infrastructure isn't prepared to guarantee that, and increasing the setting of MAXALIGN seems infeasible for multiple reasons. Adjust the code to allow use of __int128 only when we can tell the compiler to assume lesser alignment. The only known symptom of this problem so far is crashes in some parallel aggregation queries.

• Prevent stack-overflow crashes when planning extremely deeply nested set operations (UNION/INTERSECT/EXCEPT) (Tom Lane)

• Avoid crash during an EvalPlanQual recheck of an indexscan that is the inner child of a merge join (Tom Lane)

  This could only happen during an update or SELECT FOR UPDATE of a join, when there is a concurrent update of some selected row.

• Fix crash in autovacuum when extended statistics are defined for a table but can't be computed (Álvaro Herrera)

• Fix null-pointer crashes for some types of LDAP URLs appearing in pg_hba.conf (Thomas Munro)

• Prevent out-of-memory failures due to excessive growth of simple hash tables (Tomas Vondra, Andres Freund)

• Fix sample INSTR() functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)

  These functions are stated to be Oracle® compatible, but they weren't exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.

  The sample code has been adjusted to match Oracle's behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

• Fix pg_dump to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

  The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for
comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with DATABASE so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with LARGE_OBJECT. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.

Note that this change has user-visible results in the output of pg_restore --list.

- Rename pg_rewind’s copy_file_range function to avoid conflict with new Linux system call of that name (Andres Freund)

This change prevents build failures with newer glibc versions.

- In ecpg, detect indicator arrays that do not have the correct length and report an error (David Rader)

- Change the behavior of contrib/cube’s cube -> int operator to make it compatible with KNN search (Alexander Korotkov)

The meaning of the second argument (the dimension selector) has been changed to make it predictable which value is selected even when dealing with cubes of varying dimensionalities.

This is an incompatible change, but since the point of the operator was to be used in KNN searches, it seems rather useless as-is. After installing this update, any expression indexes or materialized views using this operator will need to be reindexed/refreshed.

- Avoid triggering a libc assertion in contrib/hstore, due to use of memcpy() with equal source and destination pointers (Tomas Vondra)

- Fix incorrect display of tuples’ null bitmaps in contrib/pageinspect (Maksim Milyutin)

- Fix incorrect output from contrib/pageinspect’s hash_page_items() function (Masahiko Sawada)

- In contrib/postgres_fdw, avoid “outer pathkeys do not match mergeclauses” planner error when constructing a plan involving a remote join (Robert Haas)

- In contrib/postgres_fdw, avoid planner failure when there are duplicate GROUP BY entries (Jeevan Chalke)

- Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

The scripts in contrib/start-scripts/osx use infrastructure that’s been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory contrib/start-scripts/macos containing scripts that use the newer launchd infrastructure.

- Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dunstan)

- Support linking to MinGW-built versions of libperl (Noah Misch)

This allows building PL/Perl with some common Perl distributions for Windows.

- Fix MSVC build to test whether 32-bit libperl needs -D_USE_32BIT_TIME_T (Noah Misch)

Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.

- On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)

This may allow collection of a core dump for some early-startup failures that did not produce a dump before.
Release Notes

- On Windows, avoid encoding-conversion-related crashes when emitting messages very early in postmaster startup (Takayuki Tsunakawa)
- Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)
- Add support for spinlocks on Motorola 88K (David Carlier)
- Update time zone data files to tzdata release 2018c for DST law changes in Brazil, Sao Tome and Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The US/Pacific-New zone has been removed (it was only an alias for America/Los_Angeles anyway).

E.8. Release 10.1

Release date: 2017-11-09

This release contains a variety of fixes from 10.0. For information about new features in major release 10, see Section E.9.

E.8.1. Migration to Version 10.1

A dump/restore is not required for those running 10.X.

However, if you use BRIN indexes, see the fourth changelog entry below.

E.8.2. Changes

- Ensure that INSERT ... ON CONFLICT DO UPDATE checks table permissions and RLS policies in all cases (Dean Rasheed)

  The update path of INSERT ... ON CONFLICT DO UPDATE requires SELECT permission on the columns of the arbiter index, but it failed to check for that in the case of an arbiter specified by constraint name. In addition, for a table with row level security enabled, it failed to check updated rows against the table's SELECT policies (regardless of how the arbiter index was specified). (CVE-2017-15099)

- Fix crash due to rowtype mismatch in json{b}_populate_recordset() (Michael Paquier, Tom Lane)

  These functions used the result rowtype specified in the FROM ... AS clause without checking that it matched the actual rowtype of the supplied tuple value. If it didn't, that would usually result in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)

- Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)

  Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making $PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

  By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

- Fix BRIN index summarization to handle concurrent table extension correctly (Álvaro Herrera)

  Previously, a race condition allowed some table rows to be omitted from the index. It may be necessary to reindex existing BRIN indexes to recover from past occurrences of this problem.

- Fix possible failures during concurrent updates of a BRIN index (Tom Lane)
These race conditions could result in errors like “invalid index offnum” or “inconsistent range map”.

- Prevent logical replication from setting non-replicated columns to nulls when replicating an UPDATE (Petr Jelinek)
- Fix logical replication to fire BEFORE ROW DELETE triggers when expected (Masahiko Sawada)

Previously, that failed to happen unless the table also had a BEFORE ROW UPDATE trigger.

- Fix crash when logical decoding is invoked from a SPI-using function, in particular any function written in a PL language (Tom Lane)
- Ignore CTEs when looking up the target table for INSERT/UPDATE/DELETE, and prevent matching schema-qualified target table names to trigger transition table names (Thomas Munro)

This restores the pre-v10 behavior for CTEs attached to DML commands.

- Avoid evaluating an aggregate function’s argument expression(s) at rows where its FILTER test fails (Tom Lane)

This restores the pre-v10 (and SQL-standard) behavior.

- Fix incorrect query results when multiple GROUPING SETS columns contain the same simple variable (Tom Lane)
- Fix query-lifespan memory leakage while evaluating a set-returning function in a SELECT’s target list (Tom Lane)
- Allow parallel execution of prepared statements with generic plans (Amit Kapila, Kuntal Ghosh)
- Fix incorrect parallelization decisions for nested queries (Amit Kapila, Kuntal Ghosh)
- Fix parallel query handling to not fail when a recently-used role is dropped (Amit Kapila)
- Fix crash in parallel execution of a bitmap scan having a BitmapAnd plan node below a BitmapOr node (Dilip Kumar)
- Fix JSON_BUILD_ARRAY(), JSON_BUILD_OBJECT(), and their JSONB equivalents to handle explicit VARIADIC arguments correctly (Michael Paquier)
- Fix autovacuum’s “work item” logic to prevent possible crashes and silent loss of work items (Álvaro Herrera)
- Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)
- Record proper dependencies when a view or rule contains FIELDSELECT or FIELDSTORE expression nodes (Tom Lane)

Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

- Correctly detect hashability of range data types (Tom Lane)

The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range’s subtype has hash support. This does not affect any of the built-in range types, since they’re all hashable anyway.

- Correctly ignore RELABELTYPE expression nodes when examining functional-dependency statistics (David Rowley)

This allows, e.g., extended statistics on VARCHAR columns to be used properly.
• Prevent sharing transition states between ordered-set aggregates (David Rowley)

This causes a crash with the built-in ordered-set aggregates, and probably with user-written ones as well. v11 and later will include provisions for dealing with such cases safely, but in released branches, just disable the optimization.

• Prevent `idle_in_transaction_session_timeout` from being ignored when a statement_timeout occurred earlier (Lukas Fittl)

• Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)

If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.

• Reduce the frequency of data flush requests during bulk file copies to avoid performance problems on macOS, particularly with its new APFS file system (Tom Lane)

• Allow `COPY`'s `FREEZE` option to work when the transaction isolation level is `REPEATABLE READ` or higher (Noah Misch)

This case was unintentionally broken by a previous bug fix.

• Fix `AggGetAggref()` to return the correct `Aggref` nodes to aggregate final functions whose transition calculations have been merged (Tom Lane)

• Fix insufficient schema-qualification in some new queries in `pg_dump` and `psql` (Vitaly Burovoy, Tom Lane, Noah Misch)

• Avoid use of `@>` operator in `psql`'s queries for `\d` (Tom Lane)

This prevents problems when the `parray_gin` extension is installed, since that defines a conflicting operator.

• Fix `pg_basebackup`'s matching of tablespace paths to canonicalize both paths before comparing (Michael Paquier)

This is particularly helpful on Windows.

• Fix `libpq` to not require user's home directory to exist (Tom Lane)

In v10, failure to find the home directory while trying to read `~/.pgpass` was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading `~/.pg_service.conf`, though this was less obvious since that file is not sought unless a service name is specified.

• In `ecpglib`, correctly handle backslashes in string literals depending on whether `standard_conforming_strings` is set (Tsunakawa Takayuki)

• Make `ecpglib`'s Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengqi, Michael Meskes)

• Fix missing temp-install prerequisites for `check`-like Make targets (Noah Misch)

Some non-default test procedures that are meant to work like `make check` failed to ensure that the temporary installation was up to date.

• Update time zone data files to `tzdata` release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

• In the documentation, restore HTML anchors to being upper-case strings (Peter Eisentraut)
Due to a toolchain change, the 10.0 user manual had lower-case strings for intrapage anchors, thus breaking some external links into our website documentation. Return to our previous convention of using upper-case strings.

E.9. Release 10

Release date: 2017-10-05

E.9.1. Overview

Major enhancements in PostgreSQL 10 include:

• Logical replication using publish/subscribe
• Declarative table partitioning
• Improved query parallelism
• Significant general performance improvements
• Stronger password authentication based on SCRAM-SHA-256
• Improved monitoring and control

The above items are explained in more detail in the sections below.

E.9.2. Migration to Version 10

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 10 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

• Hash indexes must be rebuilt after pg_upgrade-ing from any previous major PostgreSQL version (Mithun Cy, Robert Haas, Amit Kapila)

  Major hash index improvements necessitated this requirement. pg_upgrade will create a script to assist with this.

• Rename write-ahead log directory pg_xlog to pg_wal, and rename transaction status directory pg_clog to pg_xact (Michael Paquier)

  Users have occasionally thought that these directories contained only inessential log files, and proceeded to remove write-ahead log files or transaction status files manually, causing irrecoverable data loss. These name changes are intended to discourage such errors in future.

• Rename WAL-related functions and views to use lsn instead of location (David Rowley)

  There was previously an inconsistent mixture of the two terminologies.

• Change the implementation of set-returning functions appearing in a query’s SELECT list (Andres Freund)
Set-returning functions are now evaluated before evaluation of scalar expressions in the \texttt{SELECT} list, much as though they had been placed in a \texttt{LATERAL FROM}-clause item. This allows saner semantics for cases where multiple set-returning functions are present. If they return different numbers of rows, the shorter results are extended to match the longest result by adding nulls. Previously the results were cycled until they all terminated at the same time, producing a number of rows equal to the least common multiple of the functions' periods. In addition, set-returning functions are now disallowed within \texttt{CASE} and \texttt{COALESCE} constructs. For more information see Section 38.5.8.

- **Use standard row constructor syntax in \texttt{UPDATE ... SET \texttt{(column_list)} = row_constructor} (Tom Lane)**

  The \texttt{row_constructor} can now begin with the keyword \texttt{ROW}; previously that had to be omitted. If just one column name appears in the \texttt{column_list}, then the \texttt{row_constructor} now must use the \texttt{ROW} keyword, since otherwise it is not a valid row constructor but just a parenthesized expression. Also, an occurrence of \texttt{table_name.\ast} within the \texttt{row_constructor} is now expanded into multiple columns, as occurs in other uses of \texttt{row_constructors}.

- **When \texttt{ALTER TABLE ... ADD PRIMARY KEY} marks columns NOT NULL, that change now propagates to inheritance child tables as well** (Michael Paquier)

- **Prevent statement-level triggers from firing more than once per statement** (Tom Lane)

  Cases involving writable CTEs updating the same table updated by the containing statement, or by another writable CTE, fired \texttt{BEFORE STATEMENT} or \texttt{AFTER STATEMENT} triggers more than once. Also, if there were statement-level triggers on a table affected by a foreign key enforcement action (such as \texttt{ON DELETE CASCADE}), they could fire more than once per outer SQL statement. This is contrary to the SQL standard, so change it.

- **Move sequences' metadata fields into a new \texttt{pg_sequence} system catalog** (Peter Eisentraut)

  A sequence relation now stores only the fields that can be modified by \texttt{nextval()}, that is \texttt{last_value}, \texttt{log_cnt}, and \texttt{is_called}. Other sequence properties, such as the starting value and increment, are kept in a corresponding row of the \texttt{pg_sequence} catalog. \texttt{ALTER SEQUENCE} updates are now fully transactional, implying that the sequence is locked until commit. The \texttt{nextval()} and \texttt{setval()} functions remain nontransactional.

  The main incompatibility introduced by this change is that selecting from a sequence relation now returns only the three fields named above. To obtain the sequence's other properties, applications must look into \texttt{pg_sequence}. The new system view \texttt{pg_sequences} can also be used for this purpose; it provides column names that are more compatible with existing code.

  Also, sequences created for \texttt{SERIAL} columns now generate positive 32-bit wide values, whereas previous versions generated 64-bit wide values. This has no visible effect if the values are only stored in a column.

  The output of \texttt{psql}'s \texttt{\d} command for a sequence has been redesigned, too.

- **Make \texttt{pg_basebackup} stream the WAL needed to restore the backup by default** (Magnus Hagander)

  This changes \texttt{pg_basebackup}'s \texttt{--wal-method} default to \texttt{stream}. An option value \texttt{none} has been added to reproduce the old behavior. The \texttt{pg_basebackup} option \texttt{--x} has been removed (instead, use \texttt{--X fetch}.

- **Change how logical replication uses \texttt{pg_hba.conf}** (Peter Eisentraut)

  In previous releases, a logical replication connection required the \texttt{replication} keyword in the database column. As of this release, logical replication matches a normal entry with a database name or keywords such as \texttt{all}. Physical replication continues to use the \texttt{replication} keyword. Since built-in logical replication is new in this release, this change only affects users of third-party logical replication plugins.
• Make all pg_ctl actions wait for completion by default (Peter Eisentraut)

Previously some pg_ctl actions didn't wait for completion, and required the use of \-w to do so.

• Change the default value of the log_directory server parameter from pg_log to log (Andreas Karlsson)

• Add configuration option ssl_dh_params_file to specify file name for custom OpenSSL DH parameters (Heikki Linnakangas)

  This replaces the hardcoded, undocumented file name dh1024.pem. Note that dh1024.pem is no longer examined by default; you must set this option if you want to use custom DH parameters.

• Increase the size of the default DH parameters used for OpenSSL ephemeral DH ciphers to 2048 bits (Heikki Linnakangas)

  The size of the compiled-in DH parameters has been increased from 1024 to 2048 bits, making DH key exchange more resistant to brute-force attacks. However, some old SSL implementations, notably some revisions of Java Runtime Environment version 6, will not accept DH parameters longer than 1024 bits, and hence will not be able to connect over SSL. If it’s necessary to support such old clients, you can use custom 1024-bit DH parameters instead of the compiled-in defaults. See ssl_dh_params_file.

• Remove the ability to store unencrypted passwords on the server (Heikki Linnakangas)

  The password_encryption server parameter no longer supports off or plain. The UNENCRYPTED option is no longer supported in CREATE/ALTER USER ... PASSWORD. Similarly, the --unencrypted option has been removed from createuser. Unencrypted passwords migrated from older versions will be stored encrypted in this release. The default setting for password_encryption is still md5.

• Add min_parallel_table_scan_size and min_parallel_index_scan_size server parameters to control parallel queries (Amit Kapila, Robert Haas)

  These replace min_parallel_relation_size, which was found to be too generic.

• Don’t downcase unquoted text within shared_preload_libraries and related server parameters (QL Zhuo)

  These settings are really lists of file names, but they were previously treated as lists of SQL identifiers, which have different parsing rules.

• Remove sql_inheritance server parameter (Robert Haas)

  Changing this setting from the default value caused queries referencing parent tables to not include child tables. The SQL standard requires them to be included, however, and this has been the default since PostgreSQL 7.1.

• Allow multi-dimensional arrays to be passed into PL/Python functions, and returned as nested Python lists (Alexey Grishchenko, Dave Cramer, Heikki Linnakangas)

  This feature requires a backwards-incompatible change to the handling of arrays of composite types in PL/Python. Previously, you could return an array of composite values by writing, e.g. [[col1, col2], [col1, col2]]; but now that is interpreted as a two-dimensional array. Composite types in arrays must now be written as Python tuples, not lists, to resolve the ambiguity; that is, write [(col1, col2), (col1, col2)] instead.

• Remove PL/ Tcl’s “module” auto-loading facility (Tom Lane)

  This functionality has been replaced by new server parameters pltcl.start_proc and pltclu.start_proc, which are easier to use and more similar to features available in other PLs.
- Remove pg_dump/pg_dumpall support for dumping from pre-8.0 servers (Tom Lane)

Users needing to dump from pre-8.0 servers will need to use dump programs from PostgreSQL 9.6 or earlier. The resulting output should still load successfully into newer servers.

- Remove support for floating-point timestamps and intervals (Tom Lane)

This removes configure's --disable-integer-datetimes option. Floating-point timestamps have few advantages and have not been the default since PostgreSQL 8.3.

- Remove server support for client/server protocol version 1.0 (Tom Lane)

This protocol hasn't had client support since PostgreSQL 6.3.

- Remove contrib/tsearch2 module (Robert Haas)

This module provided compatibility with the version of full text search that shipped in pre-8.3 PostgreSQL releases.

- Remove createlang and droplang command-line applications (Peter Eisentraut)

These had been deprecated since PostgreSQL 9.1. Instead, use CREATE EXTENSION and DROP EXTENSION directly.

- Remove support for version-0 function calling conventions (Andres Freund)

Extensions providing C-coded functions must now conform to version 1 calling conventions. Version 0 has been deprecated since 2001.

### E.9.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 10 and the previous major release.

#### E.9.3.1. Server

#### E.9.3.1.1. Parallel Queries

- Support parallel B-tree index scans (Rahila Syed, Amit Kapila, Robert Haas, Rafia Sabih)

  This change allows B-tree index pages to be searched by separate parallel workers.

- Support parallel bitmap heap scans (Dilip Kumar)

  This allows a single index scan to dispatch parallel workers to process different areas of the heap.

- Allow merge joins to be performed in parallel (Dilip Kumar)

- Allow non-correlated subqueries to be run in parallel (Amit Kapila)

- Improve ability of parallel workers to return pre-sorted data (Rushabh Lathia)

- Increase parallel query usage in procedural language functions (Robert Haas, Rafia Sabih)

- Add max_parallel_workers server parameter to limit the number of worker processes that can be used for query parallelism (Julien Rouhaud)

  This parameter can be set lower than max_worker_processes to reserve worker processes for purposes other than parallel queries.

- Enable parallelism by default by changing the default setting of max_parallel_workers_per_gather to 2.
E.9.3.1.2. Indexes

- Add write-ahead logging support to hash indexes (Amit Kapila)
  This makes hash indexes crash-safe and replicatable. The former warning message about their use is removed.
- Improve hash index performance (Amit Kapila, Mithun Cy, Ashutosh Sharma)
- Add SP-GiST index support for INET and CIDR data types (Emre Hasegeli)
- Add option to allow BRIN index summarization to happen more aggressively (Álvaro Herrera)
  A new CREATE INDEX option enables auto-summarization of the previous BRIN page range when a new page range is created.
- Add functions to remove and re-add BRIN summarization for BRIN index ranges (Álvaro Herrera)
  The new SQL function brin_summarize_range() updates BRIN index summarization for a specified range and brin_desummarize_range() removes it. This is helpful to update summarization of a range that is now smaller due to UPDATEs and DELETEs.
- Improve accuracy in determining if a BRIN index scan is beneficial (David Rowley, Emre Hasegeli)
- Allow faster GiST inserts and updates by reusing index space more efficiently (Andrey Borodin)
- Reduce page locking during vacuuming of GIN indexes (Andrey Borodin)

E.9.3.1.3. Locking

- Reduce locking required to change table parameters (Simon Riggs, Fabrízio Mello)
  For example, changing a table’s effective_io_concurrency setting can now be done with a more lightweight lock.
- Allow tuning of predicate lock promotion thresholds (Dagfinn Ilmari Mannsåker)
  Lock promotion can now be controlled through two new server parameters, max_pred_locks_per_relation and max_pred_locks_per_page.

E.9.3.1.4. Optimizer

- Add multi-column optimizer statistics to compute the correlation ratio and number of distinct values (Tomas Vondra, David Rowley, Álvaro Herrera)
  New commands are CREATE STATISTICS, ALTER STATISTICS, and DROP STATISTICS. This feature is helpful in estimating query memory usage and when combining the statistics from individual columns.
- Improve performance of queries affected by row-level security restrictions (Tom Lane)
  The optimizer now has more knowledge about where it can place RLS filter conditions, allowing better plans to be generated while still enforcing the RLS conditions safely.

E.9.3.1.5. General Performance

- Speed up aggregate functions that calculate a running sum using numeric-type arithmetic, including some variants of SUM(), AVG(), and STDDEV() (Heikki Linnakangas)
- Improve performance of character encoding conversions by using radix trees (Kyotaro Horiguchi, Heikki Linnakangas)
• Reduce expression evaluation overhead during query execution, as well as plan node calling overhead (Andres Freund)

This is particularly helpful for queries that process many rows.

• Allow hashed aggregation to be used with grouping sets (Andrew Gierth)

• Use uniqueness guarantees to optimize certain join types (David Rowley)

• Improve sort performance of the macaddr data type (Brandur Leach)

• Reduce statistics tracking overhead in sessions that reference many thousands of relations (Aleksander Alekseev)

**E.9.3.1.6. Monitoring**

• Allow explicit control over EXPLAIN's display of planning and execution time (Ashutosh Bapat)

By default planning and execution time are displayed by EXPLAIN ANALYZE and are not displayed in other cases. The new EXPLAIN option SUMMARY allows explicit control of this.

• Add default monitoring roles (Dave Page)

New roles pg_monitor, pg_read_all_settings, pg_read_all_stats, and pg_stat_scan_tables allow simplified permission configuration.

• Properly update the statistics collector during REFRESH MATERIALIZED VIEW (Jim Mlodgenski)

**E.9.3.1.6.1. Logging**

• Change the default value of log_line_prefix to include current timestamp (with milliseconds) and the process ID in each line of postmaster log output (Christoph Berg)

The previous default was an empty prefix.

• Add functions to return the log and WAL directory contents (Dave Page)

The new functions are pg_ls_logdir() and pg_ls_waldir() and can be executed by non-superusers with the proper permissions.

• Add function pg_current_logfile() to read logging collector's current stderr and csvlog output file names (Gilles Darold)

• Report the address and port number of each listening socket in the server log during postmaster startup (Tom Lane)

Also, when logging failure to bind a listening socket, include the specific address we attempted to bind to.

• Reduce log chatter about the starting and stopping of launcher subprocesses (Tom Lane)

These are now DEBUG1-level messages.

• Reduce message verbosity of lower-numbered debug levels controlled by log_min_messages (Robert Haas)

This also changes the verbosity of client_min_messages debug levels.

**E.9.3.1.6.2. pg_stat_activity**

• Add pg_stat_activity reporting of low-level wait states (Michael Paquier, Robert Haas, Rushabh Lathia)
This change enables reporting of numerous low-level wait conditions, including latch waits, file reads/writes/fsyncs, client reads/writes, and synchronous replication.

• Show auxiliary processes, background workers, and walsender processes in \texttt{pg\_stat\_activity} (Kuntal Ghosh, Michael Paquier)

  This simplifies monitoring. A new column \texttt{backend\_type} identifies the process type.

• Allow \texttt{pg\_stat\_activity} to show the SQL query being executed by parallel workers (Rafia Sabih)

• Rename \texttt{pg\_stat\_activity.wait\_event\_type} values \texttt{LWLockTranche} and \texttt{LWLock-Named} to \texttt{LWLock} (Robert Haas)

  This makes the output more consistent.

\textbf{E.9.3.1.7. Authentication}

• Add SCRAM-SHA-256 support for password negotiation and storage (Michael Paquier, Heikki Linnakangas)

  This provides better security than the existing \texttt{md5} negotiation and storage method.

• Change the password\_encryption server parameter from \texttt{boolean} to \texttt{enum} (Michael Paquier)

  This was necessary to support additional password hashing options.

• Add view \texttt{pg\_hba\_file\_rules} to display the contents of \texttt{pg\_hba.conf} (Haribabu Kommi)

  This shows the file contents, not the currently active settings.

• Support multiple RADIUS servers (Magnus Hagander)

  All the RADIUS related parameters are now plural and support a comma-separated list of servers.

\textbf{E.9.3.1.8. Server Configuration}

• Allow SSL configuration to be updated during configuration reload (Andreas Karlsson, Tom Lane)

  This allows SSL to be reconfigured without a server restart, by using \texttt{pg\_ctl reload,SELECT pg\_reload\_conf()}, or sending a \texttt{SIGHUP} signal. However, reloading the SSL configuration does not work if the server's SSL key requires a passphrase, as there is no way to re-prompt for the passphrase. The original configuration will apply for the life of the postmaster in that case.

• Make the maximum value of \texttt{bgwriter\_lrui\_maxpages} effectively unlimited (Jim Nasby)

\textbf{E.9.3.1.9. Reliability}

• After creating or unlinking files, perform an fsync on their parent directory (Michael Paquier)

  This reduces the risk of data loss after a power failure.

\textbf{E.9.3.1.9.1. Write-Ahead Log (WAL)}

• Prevent unnecessary checkpoints and WAL archiving on otherwise-idle systems (Michael Paquier)

• Add \texttt{wal\_consistency\_checking} server parameter to add details to WAL that can be sanity-checked on the standby (Kuntal Ghosh, Robert Haas)

  Any sanity-check failure generates a fatal error on the standby.

• Increase the maximum configurable WAL segment size to one gigabyte (Beena Emerson)
A larger WAL segment size allows for fewer archive_command invocations and fewer WAL files to manage.

E.9.3.2. Replication and Recovery

- Add the ability to logically replicate tables to standby servers (Petr Jelinek)

Logical replication allows more flexibility than physical replication does, including replication between different major versions of PostgreSQL and selective replication.

- Allow waiting for commit acknowledgment from standby servers irrespective of the order they appear in synchronous_standby_names (Masahiko Sawada)

Previously the server always waited for the active standbys that appeared first in synchronous_standby_names. The new synchronous_standby_names keyword ANY allows waiting for any number of standbys irrespective of their ordering. This is known as quorum commit.

- Reduce configuration changes necessary to perform streaming backup and replication (Magnus Hagander, Dang Minh Huong)

Specifically, the defaults were changed for wal_level, max_wal_senders, max_replication_slots, and hot_standby to make them suitable for these usages out-of-the-box.

- Enable replication from localhost connections by default in pg_hba.conf (Michael Paquier)

Previously pg_hba.conf’s replication connection lines were commented out by default. This is particularly useful for pg_basebackup.

- Add columns to pg_stat_replication to report replication delay times (Thomas Munro)

The new columns are write_lag, flush_lag, and replay_lag.

- Allow specification of the recovery stopping point by Log Sequence Number (LSN) in recovery.conf (Michael Paquier)

Previously the stopping point could only be selected by timestamp or XID.

- Allow users to disable pg_stop_backup()’s waiting for all WAL to be archived (David Steele)

An optional second argument to pg_stop_backup() controls that behavior.

- Allow creation of temporary replication slots (Petr Jelinek)

Temporary slots are automatically removed on session exit or error.

- Improve performance of hot standby replay with better tracking of Access Exclusive locks (Simon Riggs, David Rowley)

- Speed up two-phase commit recovery performance (Stas Kelvich, Nikhil Sontakke, Michael Paquier)

E.9.3.3. Queries

- Add XMLTABLE function that converts XML-formatted data into a row set (Pavel Stehule, Álvaro Herrera)

- Fix regular expressions’ character class handling for large character codes, particularly Unicode characters above U+7FF (Tom Lane)

Previously, such characters were never recognized as belonging to locale-dependent character classes such as [:alpha:].
E.9.3.4. Utility Commands

- Add table partitioning syntax that automatically creates partition constraints and handles routing of tuple insertions and updates (Amit Langote)
  The syntax supports range and list partitioning.
- Add AFTER trigger transition tables to record changed rows (Kevin Grittner, Thomas Munro)
  Transition tables are accessible from triggers written in server-side languages.
- Allow restrictive row-level security policies (Stephen Frost)
  Previously all security policies were permissive, meaning that any matching policy allowed access. A restrictive policy must match for access to be granted. These policy types can be combined.
- When creating a foreign-key constraint, check for REFERENCES permission on only the referenced table (Tom Lane)
  Previously REFERENCES permission on the referencing table was also required. This appears to have stemmed from a misreading of the SQL standard. Since creating a foreign key (or any other type of) constraint requires ownership privilege on the constrained table, additionally requiring REFERENCES permission seems rather pointless.
- Allow default permissions on schemas (Matheus Oliveira)
  This is done using the ALTER DEFAULT PRIVILEGES command.
- Add CREATE SEQUENCE AS command to create a sequence matching an integer data type (Peter Eisentraut)
  This simplifies the creation of sequences matching the range of base columns.
- Allow COPY view FROM source on views with INSTEAD INSERT triggers (Haribabu Kommi)
  The triggers are fed the data rows read by COPY.
- Allow the specification of a function name without arguments in DDL commands, if it is unique (Peter Eisentraut)
  For example, allow DROP FUNCTION on a function name without arguments if there is only one function with that name. This behavior is required by the SQL standard.
- Allow multiple functions, operators, and aggregates to be dropped with a single DROP command (Peter Eisentraut)
- Support IF NOT EXISTS in CREATE SERVER, CREATE USER MAPPING, and CREATE COLLATION (Anastasia Lubennikova, Peter Eisentraut)
- Make VACUUM VERBOSE report the number of skipped frozen pages and oldest xmin (Masahiko Sawada, Simon Riggs)
  This information is also included in log_autovacuum_min_duration output.
- Improve speed of VACUUM's removal of trailing empty heap pages (Claudio Freire, Álvaro Herrera)

E.9.3.5. Data Types

- Add full text search support for JSON and JSONB (Dmitry Dolgov)
  The functions ts_headline() and to_tsvector() can now be used on these data types.
• Add support for EUI-64 MAC addresses, as a new data type `macaddr8` (Haribabu Kommi)

  This complements the existing support for EUI-48 MAC addresses (type `macaddr`).

• Add identity columns for assigning a numeric value to columns on insert (Peter Eisentraut)

  These are similar to `SERIAL` columns, but are SQL standard compliant.

• Allow `ENUM` values to be renamed (Dagfinn Ilmari Mannsåker)

  This uses the syntax `ALTER TYPE ... RENAME VALUE`.

• Properly treat array pseudotypes (`anyarray`) as arrays in `to_json()` and `to_jsonb()` (Andrew Dunstan)

  Previously columns declared as `anyarray` (particularly those in the `pg_stats` view) were converted to JSON strings rather than arrays.

• Add operators for multiplication and division of `money` values with `int8` values (Peter Eisentraut)

  Previously such cases would result in converting the `int8` values to `float8` and then using the `money-and-float8` operators. The new behavior avoids possible precision loss. But note that division of `money` by `int8` now truncates the quotient, like other integer-division cases, while the previous behavior would have rounded.

• Check for overflow in the `money` type’s input function (Peter Eisentraut)

E.9.3.6. Functions

• Add simplified `regexp_match()` function (Emre Hasegeli)

  This is similar to `regexp_matches()`, but it only returns results from the first match so it does not need to return a set, making it easier to use for simple cases.

• Add a version of `jsonb`’s delete operator that takes an array of keys to delete (Magnus Hagander)

• Make `json_populate_record()` and related functions process JSON arrays and objects recursively (Nikita Glukhov)

  With this change, array-type fields in the destination SQL type are properly converted from JSON arrays, and composite-type fields are properly converted from JSON objects. Previously, such cases would fail because the text representation of the JSON value would be fed to `array_in()` or `record_in()`, and its syntax would not match what those input functions expect.

• Add function `txid_current_if_assigned()` to return the current transaction ID or NULL if no transaction ID has been assigned (Craig Ringer)

  This is different from `txid_current()`, which always returns a transaction ID, assigning one if necessary. Unlike that function, this function can be run on standby servers.

• Add function `txid_status()` to check if a transaction was committed (Craig Ringer)

  This is useful for checking after an abrupt disconnection whether your previous transaction committed and you just didn’t receive the acknowledgment.

• Allow `make_date()` to interpret negative years as BC years (Álvaro Herrera)

• Make `to_timestamp()` and `to_date()` reject out-of-range input fields (Artur Zakirov)

  For example, previously `to_date('2009-06-40','YYYY-MM-DD')` was accepted and returned `2009-07-10`. It will now generate an error.
E.9.3.7. Server-Side Languages

- Allow PL/Python's `cursor()` and `execute()` functions to be called as methods of their plan-object arguments (Peter Eisentraut)
  
  This allows a more object-oriented programming style.

- Allow PL/pgSQL's `GET DIAGNOSTICS` statement to retrieve values into array elements (Tom Lane)
  
  Previously, a syntactic restriction prevented the target variable from being an array element.

E.9.3.7.1. PL/Tcl

- Allow PL/Tcl functions to return composite types and sets (Karl Lehenbauer)

- Add a subtransaction command to PL/Tcl (Victor Wagner)
  
  This allows PL/Tcl queries to fail without aborting the entire function.

- Add server parameters `pltcl.start_proc` and `pltclu.start_proc`, to allow initialization functions to be called on PL/Tcl startup (Tom Lane)

E.9.3.8. Client Interfaces

- Allow specification of multiple host names or addresses in libpq connection strings and URIs (Robert Haas, Heikki Linnakangas)
  
  libpq will connect to the first responsive server in the list.

- Allow libpq connection strings and URIs to request a read/write host, that is a master server rather than a standby server (Victor Wagner, Mithun Cy)
  
  This is useful when multiple host names are specified. It is controlled by libpq connection parameter `target_session_attrs`.

- Allow the password file name to be specified as a libpq connection parameter (Julian Markwort)
  
  Previously this could only be specified via an environment variable.

- Add function `PQencryptPasswordConn()` to allow creation of more types of encrypted passwords on the client side (Michael Paquier, Heikki Linnakangas)
  
  Previously only MD5-encrypted passwords could be created using `PQencryptPassword()`. This new function can also create `SCRAM-SHA-256`-encrypted passwords.

- Change ecpg preprocessor version from 4.12 to 10 (Tom Lane)
  
  Henceforth the ecpg version will match the PostgreSQL distribution version number.

E.9.3.9. Client Applications

E.9.3.9.1. psql

- Add conditional branch support to psql (Corey Huinker)
  
  This feature adds psql meta-commands `\if, \elif, \else, and \endif`. This is primarily helpful for scripting.

- Add psql `\gx` meta-command to execute (`\g`) a query in expanded mode (`\x`) (Christoph Berg)
• Expand \texttt{psql} variable references in backtick-executed strings (Tom Lane)

This is particularly useful in the new \texttt{psql} conditional branch commands.

• Prevent \texttt{psql}'s special variables from being set to invalid values (Daniel Vérité, Tom Lane)

Previously, setting one of \texttt{psql}'s special variables to an invalid value silently resulted in the default behavior, \texttt{\textbackslash set} on a special variable now fails if the proposed new value is invalid. As a special exception, \texttt{\textbackslash set} with an empty or omitted new value, on a boolean-valued special variable, still has the effect of setting the variable to \texttt{on}; but now it actually acquires that value rather than an empty string. \texttt{\textbackslash unset} on a special variable now explicitly sets the variable to its default value, which is also the value it acquires at startup. In sum, a control variable now always has a displayable value that reflects what \texttt{psql} is actually doing.

• Add variables showing server version and \texttt{psql} version (Fabien Coelho)

• Improve \texttt{psql}'s \texttt{\textbackslash d} (display relation) and \texttt{\textbackslash dD} (display domain) commands to show collation, nullable, and default properties in separate columns (Peter Eisentraut)

Previously they were shown in a single “Modifiers” column.

• Make the various \texttt{\textbackslash d} commands handle no-matching-object cases more consistently (Daniel Gustafsson)

They now all print the message about that to stderr, not stdout, and the message wording is more consistent.

• Improve \texttt{psql}'s tab completion (Jeff Janes, Ian Barwick, Andreas Karlsson, Sehope Sarkuni, Thomas Munro, Kevin Grittner, Dagfinn Ilmari Mannsåker)

E.9.3.9.2. pgbench

• Add pgbench option --log-prefix to control the log file prefix (Masahiko Sawada)

• Allow pgbench’s meta-commands to span multiple lines (Fabien Coelho)

A meta-command can now be continued onto the next line by writing backslash-return.

• Remove restriction on placement of --M option relative to other command line options (Tom Lane)

E.9.3.10. Server Applications

• Add pg_receivewal option --Z/--compress to specify compression (Michael Paquier)

• Add pg_recvlogical option --endpos to specify the ending position (Craig Ringer)

This complements the existing --startpos option.

• Rename initdb options --noclean and --nosync to be spelled --no-clean and --no-sync (Vik Fearing, Peter Eisentraut)

The old spellings are still supported.

E.9.3.10.1. pg_dump, pg_dumpall, pg_restore

• Allow pg_restore to exclude schemas (Michael Banck)

This adds a new --N/--exclude-schema option.

• Add --no-blobs option to pg_dump (Guillaume Lelarge)

This suppresses dumping of large objects.
Release Notes

- Add `pg_dumpall` option `--no-role-passwords` to omit role passwords (Robins Tharakan, Simon Riggs)

  This allows use of `pg_dumpall` by non-superusers; without this option, it fails due to inability to read passwords.

- Support using synchronized snapshots when dumping from a standby server (Petr Jelinek)

- Issue `fsync()` on the output files generated by `pg_dump` and `pg_dumpall` (Michael Paquier)

  This provides more security that the output is safely stored on disk before the program exits. This can be disabled with the new `--no-sync` option.

E.9.3.10.2. pg_basebackup

- Allow `pg_basebackup` to stream write-ahead log in tar mode (Magnus Hagander)

  The WAL will be stored in a separate tar file from the base backup.

- Make `pg_basebackup` use temporary replication slots (Magnus Hagander)

  Temporary replication slots will be used by default when `pg_basebackup` uses WAL streaming with default options.

- Be more careful about `fsync`'ing in all required places in `pg_basebackup` and `pg_receivewal` (Michael Paquier)

- Add `pg_basebackup` option `--no-sync` to disable `fsync` (Michael Paquier)

- Improve `pg_basebackup`'s handling of which directories to skip (David Steele)

E.9.3.10.3. pg_ctl

- Add wait option for `pg_ctl`'s promote operation (Peter Eisentraut)

- Add long options for `pg_ctl` wait (`--wait`) and no-wait (`--no-wait`) (Vik Fearing)

- Add long option for `pg_ctl` server options (`--options`) (Peter Eisentraut)

- Make `pg_ctl` start `--wait` detect server-ready by watching `postmaster.pid`, not by attempting connections (Tom Lane)

  The postmaster has been changed to report its ready-for-connections status in `postmaster.pid`, and `pg_ctl` now examines that file to detect whether startup is complete. This is more efficient and reliable than the old method, and it eliminates postmaster log entries about rejected connection attempts during startup.

- Reduce `pg_ctl`'s reaction time when waiting for postmaster start/stop (Tom Lane)

  `pg_ctl` now probes ten times per second when waiting for a postmaster state change, rather than once per second.

- Ensure that `pg_ctl` exits with nonzero status if an operation being waited for does not complete within the timeout (Peter Eisentraut)

  The `start` and `promote` operations now return exit status 1, not 0, in such cases. The `stop` operation has always done that.

E.9.3.11. Source Code

- Change to two-part release version numbering (Peter Eisentraut, Tom Lane)
Release numbers will now have two parts (e.g., 10.1) rather than three (e.g., 9.6.3). Major versions will now increase just the first number, and minor releases will increase just the second number. Release branches will be referred to by single numbers (e.g., 10 rather than 9.6). This change is intended to reduce user confusion about what is a major or minor release of PostgreSQL.

- Improve behavior of pgindent (Piotr Stefaniak, Tom Lane)

We have switched to a new version of pg_bsd_indent based on recent improvements made by the FreeBSD project. This fixes numerous small bugs that led to odd C code formatting decisions. Most notably, lines within parentheses (such as in a multi-line function call) are now uniformly indented to match the opening paren, even if that would result in code extending past the right margin.

- Allow the ICU library to optionally be used for collation support (Peter Eisentraut)

The ICU library has versioning that allows detection of collation changes between versions. It is enabled via configure option --with-icu. The default still uses the operating system's native collation library.

- Automatically mark all PG_FUNCTION_INFO_V1 functions as DLLEXPORT-ed on Windows (Laurenz Albe)

If third-party code is using extern function declarations, they should also add DLLEXPORT markers to those declarations.

- Remove SPI functions SPI_push(), SPI_pop(), SPI_push_conditional(), SPI_pop_conditional(), and SPI_restore_connection() as unnecessary (Tom Lane)

Their functionality now happens automatically. There are now no-op macros by these names so that external modules don't need to be updated immediately, but eventually such calls should be removed.

A side effect of this change is that SPI_palloc() and allied functions now require an active SPI connection; they do not degenerate to simple palloc() if there is none. That previous behavior was not very useful and posed risks of unexpected memory leaks.

- Allow shared memory to be dynamically allocated (Thomas Munro, Robert Haas)

- Add slab-like memory allocator for efficient fixed-size allocations (Tomas Vondra)

- Use POSIX semaphores rather than SysV semaphores on Linux and FreeBSD (Tom Lane)

This avoids platform-specific limits on SysV semaphore usage.

- Improve support for 64-bit atomics (Andres Freund)

- Enable 64-bit atomic operations on ARM64 (Roman Shaposhnik)

- Switch to using clock_gettime(), if available, for duration measurements (Tom Lane)

  gettimeofday() is still used if clock_gettime() is not available.

- Add more robust random number generators to be used for cryptographically secure uses (Magnus Hagander, Michael Paquier, Heikki Linnakangas)

  If no strong random number generator can be found, configure will fail unless the --disable-strong-random option is used. However, with this option, pgcrypto functions requiring a strong random number generator will be disabled.

- Allow WaitLatchOrSocket() to wait for socket connection on Windows (Andres Freund)
• **tupconvert.c** functions no longer convert tuples just to embed a different composite-type OID in them (Ashutosh Bapat, Tom Lane)

  The majority of callers don't care about the composite-type OID; but if the result tuple is to be used as a composite Datum, steps should be taken to make sure the correct OID is inserted in it.

• Remove SCO and Unixware ports (Tom Lane)

• Overhaul documentation build process (Alexander Lakhin)

• Use XSLT to build the PostgreSQL documentation (Peter Eisentraut)

  Previously Jade, DSSSL, and JadeTex were used.

• Build HTML documentation using XSLT stylesheets by default (Peter Eisentraut)

### E.9.3.12. Additional Modules

• Allow file_fdw to read from program output as well as files (Corey Huinker, Adam Gomaa)

• In postgres_fdw, push aggregate functions to the remote server, when possible (Jeevan Chalke, Ashutosh Bapat)

  This reduces the amount of data that must be passed from the remote server, and offloads aggregate computation from the requesting server.

• In postgres_fdw, push joins to the remote server in more cases (David Rowley, Ashutosh Bapat, Etsuro Fujita)

• Properly support OID columns in postgres_fdw tables (Etsuro Fujita)

  Previously OID columns always returned zeros.

• Allow btree_gist and btree_gin to index enum types (Andrew Dunstan)

  This allows enums to be used in exclusion constraints.

• Add indexing support to btree_gist for the UUID data type (Paul Jungwirth)

• Add amcheck which can check the validity of B-tree indexes (Peter Geoghegan)

• Show ignored constants as $N rather than ? in pg_stat_statements (Lukas Fittl)

• Improve cube's handling of zero-dimensional cubes (Tom Lane)

  This also improves handling of infinite and NaN values.

• Allow pg_buffercache to run with fewer locks (Ivan Kartyshov)

  This makes it less disruptive when run on production systems.

• Add pgstattuple function pgstathashindex() to view hash index statistics (Ashutosh Sharma)

• Use GRANT permissions to control pgstattuple function usage (Stephen Frost)

  This allows DBAs to allow non-superusers to run these functions.

• Reduce locking when pgstattuple examines hash indexes (Amit Kapila)

• Add pageinspect function page_checksum() to show a page's checksum (Tomas Vondra)

• Add pageinspect function bt_page_items() to print page items from a page image (Tomas Vondra)

• Add hash index support to pageinspect (Jesper Pedersen, Ashutosh Sharma)
E.9.4. Acknowledgments

The following individuals (in alphabetical order) have contributed to this release as patch authors, committers, reviewers, testers, or reporters of issues.

Adam Brightwell
Adam Brusselback
Adam Gomaa
Adam Sah
Adrian Klaver
Aidan Van Dyk
Aleksander Alekseev
Alexander Korotkov
Alexander Lakhin
Alexander Sosna
Alexey Bashtanov
Alexey Grishchenko
Alexey Isayko
Álvaro Hernández Tortosa
Álvaro Herrera
Amit Kapila
Amit Khandekar
Amit Langote
Amul Sul
Anastasia Lubennikova
Andreas Joseph Krogh
Andreas Karlsson
Andreas Scherbaum
Andreas Seltenreich
Andres Freund
Andrew Dunstan
Andrew Gierth
Andrew Wheelwright
Andrey Borodin
Andrey Lizenko
Andy Abelisto
Antonin Houska
Arts Aasma
Arjen Nienhuis
Arseny Sher
Artur Zakirov
Ashutosh Bapat
Ashutosh Sharma
Ashwin Agrawal
Atsushi Torikoshi
Ayumi Ishii
Basil Bourque
Beena Emerson
Ben de Graaff
Benedikt Grundmann
Bernd Helmle
Brad DeJong
Brandur Leach
Breen Hagan
Bruce Momjian
Bruno Wolff III
Catalin Iacob
Chapman Flack
Chen Huajun
Choi Doo-Won
Chris Bandy
Chris Richards
Chris Ruprecht
Christian Ullrich
Christoph Berg
Chuanting Wang
Claudio Freire
Clinton Adams
Const Zhang
Constantin Pan
Corey Huinker
Craig Ringer
Cynthia Shang
Dagfinn Ilmari Mannsåker
Daisuke Higuchi
Damian Quiroga
Dan Wood
Dang Minh Huong
Daniel Gustafsson
Daniel Vérité
Daniel Westermann
Daniele Varrazzo
Danylo Hlynskyi
Darko Prelec
Dave Cramer
Dave Page
David Christensen
David Fetter
David Johnston
David Rader
David Rowley
David Steele
Dean Rasheed
Denis Smirnov
Denish Patel
Dennis Björklund
Devrim Gündüz
Dilip Kumar
Dilyan Palauzov
Dima Pavlov
Dimitry Ivanov
Dmitriy Sarafannikov
Dmitry Dolgov
Dmitry Fedin
Don Morrison
Egor Rogov
Eiji Seki
Emil Iggland
Emre Hasegeli
Enrique Meneses
Erik Nordström
Erik Rijkers
Erwin Brandstetter
Etsuro Fujita
Eugen Konkov
Eugene Kazakov
Euler Taveira
Fabien Coelho
Fabrício de Royes Mello
Feike Steenbergen
Felix Gerzaguet
Filip Jirsák
Fujii Masao
Gabriele Bartolini
Gabrielle Roth
Gao Zengqi
Gerdan Santos
Gianni Ciolli
Gilles Darold
Giuseppe Broccolo
Graham Dutton
Greg Atkins
Greg Burek
Grigory Smolkin
Guillaume Lelarge
Hans Buschmann
Haribabu Kommi
Heikki Linnakangas
Henry Boehlert
Huan Ruan
Ian Barwick
Igor Korot
Ildus Kurbangaliev
Ivan Kartyshov
Jaime Casanova
Jakob Egger
James Parks
Jarred Ward
Jason Li
Jason O'Donnell
Jason Petersen
Jeevan Chalke
Jeevan Ladhe
Jeff Dafoe
Jeff Davis
Jeff Janes
Jelte Fennema
Jeremy Finzel
Jeremy Schneider
Jeroen van der Ham
Jesper Pedersen
Jim Mlodgenski
Jim Nasby
Jinyu Zhang
Joe Conway
Joel Jacobson
John Harvey
Jon Nelson
Jordan Gigov
Josh Berkus
Josh Soref
Julian Markwort
Julien Rouhaud
Junseok Yang
E.10. Release 9.6.11

Release date: 2018-11-08

This release contains a variety of fixes from 9.6.10. For information about new features in the 9.6 major release, see Section E.21.

E.10.1. Migration to Version 9.6.11

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.12.

E.10.2. Changes

- Fix corner-case failures in has_foo_privilege() family of functions (Tom Lane)
Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. has_column_privilege() was additionally capable of crashing on some platforms.

- Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)
- Fix parsing of standard multi-character operators that are immediately followed by a comment or + or − (Andrew Gierth)

This oversight could lead to parse errors, or to incorrect assignment of precedence.

- Avoid O(N^3) slowdown in lexer for long strings of + or − characters (Andrew Gierth)
- Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)
- Fix failure of UPDATE/DELETE ... WHERE CURRENT OF ... after rewinding the referenced cursor (Tom Lane)

A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

- Fix EvalPlanQual to handle conditionally-executed InitPlans properly (Andrew Gierth, Tom Lane)
- Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)
- Disallow pushing sub-SELECTs containing window functions, LIMIT, or OFFSET to parallel workers (Amit Kapila)

Such cases could result in inconsistent behavior due to different workers getting different answers, as a result of indeterminacy due to row-ordering variations.

- Ensure that sequences owned by a foreign table are processed by ALTER OWNER on the table (Peter Eisentraut)
- Ensure that the server will process already-received NOTIFY and SIGTERM interrupts before waiting for client input (Jeff Janes, Tom Lane)
- Fix over-allocation of space for array_out()’s result string (Keiichi Hirobe)
- Fix memory leak in repeated SP-GiST index scans (Tom Lane)

This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- Ensure that ApplyLogicalMappingFile() closes the mapping file when done with it (Tomas Vondra)
  Previously, the file descriptor was leaked, eventually resulting in failures during logical decoding.
- Fix logical decoding to handle cases where a mapped catalog table is repeatedly rewritten, e.g. by VACUUM FULL (Andres Freund)
Release Notes

- Prevent starting the server with `wal_level` set to too low a value to support an existing replication slot (Andres Freund)

- Avoid crash if a utility command causes infinite recursion (Tom Lane)

- When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master (Michael Paquier, Konstantin Knizhnik)

- Fix event triggers to handle nested `ALTER TABLE` commands (Michael Paquier, Álvaro Herrera)

- Propagate parent process's transaction and statement start timestamps to parallel workers (Konstantin Knizhnik)

  This prevents misbehavior of functions such as `transaction_timestamp()` when executed in a worker.

- Fix transfer of expanded datums to parallel workers so that alignment is preserved, preventing crashes on alignment-picky platforms (Tom Lane, Amit Kapila)

- Fix WAL file recycling logic to work correctly on standby servers (Michael Paquier)

  Depending on the setting of `archive_mode`, a standby might fail to remove some WAL files that could be removed.

- Fix handling of commit-timestamp tracking during recovery (Masahiko Sawasa, Michael Paquier)

  If commit timestamp tracking has been turned on or off, recovery might fail due to trying to fetch the commit timestamp for a transaction that did not record it.

- Randomize the `random()` seed in bootstrap and standalone backends, and in `initdb` (Noah Misch)

  The main practical effect of this change is that it avoids a scenario where `initdb` might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

- Allow DSM allocation to be interrupted (Chris Travers)

- Avoid failure in a parallel worker when loading an extension that tries to access system caches within its init function (Thomas Munro)

  We don't consider that to be good extension coding practice, but it mostly worked before parallel query, so continue to support it for now.

- Properly handle turning `full_page_writes` on dynamically (Kyotaro Horiguchi)

- Fix possible crash due to double `free()` during SP-GiST rescan (Andrew Gierth)

- Avoid possible buffer overrun when replaying GIN page recompression from WAL (Alexander Korotkov, Sivasubramanian Ramasubramanian)

- Fix missed `fsync` of a replication slot's directory (Konstantin Knizhnik, Michael Paquier)

- Fix unexpected timeouts when using `wal_sender_timeout` on a slow server (Noah Misch)

- Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)

  This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.

- Ensure background workers are stopped properly when the postmaster receives a fast-shutdown request before completing database startup (Alexander Kukushkin)

2425
- Update the free space map during WAL replay of page all-visible/frozen flag changes (Álvaro Herrera)

Previously we were not careful about this, reasoning that the FSM is not critical data anyway. However, if it's sufficiently out of date, that can result in significant performance degradation after a standby has been promoted to primary. The FSM will eventually be healed by updates, but we'd like it to be good sooner, so work harder at maintaining it during WAL replay.

- Avoid premature release of parallel-query resources when query end or tuple count limit is reached (Amit Kapila)

It's only okay to shut down the executor at this point if the caller cannot demand backwards scan afterwards.

- Don't run atexit callbacks when servicing SIGQUIT (Heikki Linnakangas)

- Don't record foreign-server user mappings as members of extensions (Tom Lane)

If CREATE USER MAPPING is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can't be extension members, so user mappings shouldn't be either.

- Make syslogger more robust against failures in opening CSV log files (Tom Lane)

- Fix psql, as well as documentation examples, to call PQconsumeInput() before each PQnotifies() call (Tom Lane)

This fixes cases in which psql would not report receipt of a NOTIFY message until after the next command.

- Fix possible inconsistency in pg_dump's sorting of dissimilar object names (Jacob Champion)

- Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)

This avoids failures due to the new policy of running restores with restrictive search path.

- Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)

pg_upgrade failed to preserve an event trigger's extension-membership status.

- Fix pg_upgrade's cluster state check to work correctly on a standby server (Bruce Momjijian)

- Enforce type cube's dimension limit in all contrib/cube functions (Andrey Borodin)

Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.

- In contrib/postgres_fdw, don't try to ship a variable-free ORDER BY clause to the remote server (Andrew Gierth)

- Fix contrib/unaccent's unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)

Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.

- Fix build problems on macOS 10.14 (Mojave) (Tom Lane)

Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.
It is now recommended that Perl-related extensions write \$(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.

Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.

• Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)

Perl no longer includes the current directory in its search path by default; work around that.

• On Windows, allow the regression tests to be run by an Administrator account (Andrew Dunstan)

To do this safely, pg_regress now gives up any such privileges at startup.

• Allow btree comparison functions to return INT_MIN (Tom Lane)

Up to now, we’ve forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of \(\text{memcmp}()\), \(\text{strcmp}()\), etc, as POSIX doesn’t place any such restriction on those functions. At least some recent versions of \(\text{memcmp}()\) can return INT_MIN, causing incorrect sort ordering. Hence, we’ve removed this restriction. Callers must now use the \(\text{INVERT_COMPARE_RESULT}()\) macro if they wish to invert the sort order.

• Fix recursion hazard in shared-invalidation message processing (Tom Lane)

This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.

This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.

• Save and restore SPI’s global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)

This prevents possible interference when one SPI-using function calls another.

• Avoid using potentially-under-aligned page buffers (Tom Lane)

Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can’t place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.

• Make src/port/snprintf.c follow the C99 standard’s definition of \text{snprintf}()’s result value (Tom Lane)

On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.

• When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)

This avoids problems with missed floating point overflow checks.

• Fix configure’s detection of the result type of \text{strerror_r}() (Tom Lane)

The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.
• Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.11. Release 9.6.10

Release date: 2018-08-09

This release contains a variety of fixes from 9.6.9. For information about new features in the 9.6 major release, see Section E.21.

E.11.1. Migration to Version 9.6.10

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.12.

E.11.2. Changes

• Fix failure to reset libpq’s state fully between connection attempts (Tom Lane)

An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general, an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasichkov for reporting this issue. (CVE-2018-10915)

• Fix INSERT ... ON CONFLICT UPDATE through a view that isn't just SELECT * FROM ...

(Dean Rasheed, Amit Langote)

Erroneous expansion of an updatable view could lead to crashes or “attribute ... has the wrong type” errors, if the view's SELECT list doesn't match one-to-one with the underlying table's columns. Furthermore, this bug could be leveraged to allow updates of columns that an attacking user lacks UPDATE privilege for, if that user has INSERT and UPDATE privileges for some other column(s) of the table. Any user could also use it for disclosure of server memory. (CVE-2018-10925)

• Ensure that updates to the relfrozenxid and relminmxid values for “nailed” system catalogs are processed in a timely fashion (Andres Freund)

Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as pg_authid, because the stale cache data could persist into new sessions as well as existing ones.

• Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

• Avoid emitting a bogus WAL record when recycling an all-zero btree page (Amit Kapila)

This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

• During WAL replay, guard against corrupted record lengths exceeding 1GB (Michael Paquier)
Treat such a case as corrupt data. Previously, the code would try to allocate space and get a hard error, making recovery impossible.

- When ending recovery, delay writing the timeline history file as long as possible (Heikki Linnakan-gas)

  This avoids some situations where a failure during recovery cleanup (such as a problem with a two-phase state file) led to inconsistent timeline state on-disk.

- Improve performance of WAL replay for transactions that drop many relations (Fujii Masao)

  This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

- Improve performance of lock releasing in standby server WAL replay (Thomas Munro)

- Make logical WAL senders report streaming state correctly (Simon Riggs, Sawada Masahiko)

  The code previously mis-detected whether or not it had caught up with the upstream server.

- Fix bugs in snapshot handling during logical decoding, allowing wrong decoding results in rare cases (Arseny Sher, Álvaro Herrera)

- Ensure a table's cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)

  Previously, the failed index's OID could remain in the list, causing problems later in the same session.

- Fix mishandling of empty uncompressed posting list pages in GIN indexes (Sivasubramanian Ra-masubramanian, Alexander Korotkov)

  This could result in an assertion failure after pg_upgrade of a pre-9.4 GIN index (9.4 and later will not create such pages).

- Ensure that VACUUM will respond to signals within btree page deletion loops (Andres Freund)

  Corrupted btree indexes could result in an infinite loop here, and that previously wasn't interruptible without forcing a crash.

- Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)

  This resulted in failure to recognize that an index on a composite column could provide the sort order needed for a mergejoin on that column.

- Fix planner to avoid “ORDER/GROUP BY expression not found in targetlist” errors in some queries with set-returning functions (Tom Lane)

- Fix SQL-standard FETCH FIRST syntax to allow parameters (\$n), as the standard expects (An-drew Gierth)

- Fix EXPLAIN's accounting for resource usage, particularly buffer accesses, in parallel workers (Amit Kapila, Robert Haas)

- Fix failure to schema-qualify some object names in getObjectDescription output (Kyotaro Horiguchi, Tom Lane)

  Names of collations, conversions, and text search objects were not schema-qualified when they should be.

- Fix CREATE AGGREGATE type checking so that parallelism support functions can be attached to variadic aggregates (Alexey Bashtanov)
• Widen COPY FROM's current-line-number counter from 32 to 64 bits (David Rowley)

This avoids two problems with input exceeding 4G lines: COPY FROM WITH HEADER would drop a line every 4G lines, not only the first line, and error reports could show a wrong line number.

• Add a string freeing function to ecpg's pgtypes library, so that cross-module memory management problems can be avoided on Windows (Takayuki Tsunakawa)

On Windows, crashes can ensue if the free call for a given chunk of memory is not made from the same DLL that malloc'ed the memory. The pgtypes library sometimes returns strings that it expects the caller to free, making it impossible to follow this rule. Add a PGTYPE_Schar_free() function that just wraps free, allowing applications to follow this rule.

• Fix ecpg's support for long long variables on Windows, as well as other platforms that declare strtoll/strtoul nonstandardly or not at all (Dang Minh Huong, Tom Lane)

• Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in the semantics of a statement intra-session (Tom Lane)

This error led to assertion failures, or in rare cases, failure to enforce the INTO STRICT option as expected.

• Fix password prompting in client programs so that echo is properly disabled on Windows when stdin is not the terminal (Matthew Stickney)

• Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)

The previous fix for quoting of search_path and other list-valued variables in pg_dump output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

• Fix pg_dump's failure to dump REPLICA IDENTITY properties for constraint indexes (Tom Lane)

Manually created unique indexes were properly marked, but not those created by declaring UNIQUE or PRIMARY KEY constraints.

• Make pg_upgrade check that the old server was shut down cleanly (Bruce Momjian)

The previous check could be fooled by an immediate-mode shutdown.

• Fix contrib/hstore_plperl to look through Perl scalar references, and to not crash if it doesn't find a hash reference where it expects one (Tom Lane)

• Fix crash in contrib/ltree's lca() function when the input array is empty (Pierre Ducroquet)

• Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

• Rearrange makefiles to ensure that programs link to freshly-built libraries (such as libpq.so) rather than ones that might exist in the system library directories (Tom Lane)

This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

• Update time zone data files to tzdata release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.

This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbreviations in use; the only known effect is that the is_dst column in the pg_timezone_names view will now be true in winter and false in summer in these cases.
E.12. Release 9.6.9

Release date: 2018-05-10

This release contains a variety of fixes from 9.6.8. For information about new features in the 9.6 major release, see Section E.21.

E.12.1. Migration to Version 9.6.9

A dump/restore is not required for those running 9.6.X.

However, if you use the adminpack extension, you should update it as per the first changelog entry below.

Also, if the function marking mistakes mentioned in the second and third changelog entries below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 9.6.8, see Section E.13.

E.12.2. Changes

• Remove public execute privilege from contrib/adminpack's pg_logfile_rotate() function (Stephen Frost)

pg_logfile_rotate() is a deprecated wrapper for the core function pg_rotate_logfile(). When that function was changed to rely on SQL privileges for access control rather than a hard-coded superuser check, pg_logfile_rotate() should have been updated as well, but the need for this was missed. Hence, if adminpack is installed, any user could request a logfile rotation, creating a minor security issue.

After installing this update, administrators should update adminpack by performing ALTER EXTENSION adminpack UPDATE in each database in which adminpack is installed. (CVE-2018-1115)

• Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

The functions query_to_xml, cursor_to_xml, cursor_to_xmlschema, query_to_xmlschema, and query_to_xml and xmlschema should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions' pg_proc entries, for example ALTER FUNCTION pg_catalog.query_to_xml(text, boolean, boolean, text) VOLATILE. (Note that that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

• Fix incorrect parallel-safety markings on a few built-in functions (Thomas Munro, Tom Lane)

The functions brin_summarize_new_values, gin_clean_pending_list, cursor_to_xml, cursor_to_xmlschema, ts_rewrite, ts_stat, and binary_upgrade_create_empty_extension should be marked parallel-unsafe; some because they perform database modifications directly, and others because they execute user-supplied queries that might do so. They were marked parallel-restricted instead, leading to a risk of unexpected query errors. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard unless force_parallel_mode is turned on. In case of trouble, it can be fixed by manually updating these functions' pg_proc entries, for example AL-
TER FUNCTION pg_catalog.brin_summarize_new_values(regclass) PARALLEL UNSAFE. (Note that that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

- Avoid re-using TOAST value OIDs that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

Once the OID counter has wrapped around, it's possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value nnnnn” errors, which would persist until the dead entry was removed by VACUUM. Fix by not selecting such OIDs when creating a new TOAST entry.

- Change ANALYZE's algorithm for updating pg_class.reltuples (David Gould)

Previously, pages not actually scanned by ANALYZE were assumed to retain their old tuple density. In a large table where ANALYZE samples only a small fraction of the pages, this meant that the overall tuple density estimate could not change very much, so that reltuples would change nearly proportionally to changes in the table's physical size (retpages) regardless of what was actually happening in the table. This has been observed to result in reltuples becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that ANALYZE's sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

- Avoid deadlocks in concurrent CREATE INDEX CONCURRENTLY commands that are run under SERIALIZABLE or REPEATABLE TRANSACTIONS isolation (Tom Lane)

- Fix possible slow execution of REFRESH MATERIALIZED VIEW CONCURRENTLY (Thomas Munro)

- Fix UPDATE/DELETE ... WHERE CURRENT OF to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)

- Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)

This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.

- Fix possibly incorrect generation of an index-only-scan plan when the same table column appears in multiple index columns, and only some of those index columns use operator classes that can return the column value (Kyotaro Horiguchi)

- Fix misoptimization of CHECK constraints having provably-NULL subclauses of top-level AND/OR conditions (Tom Lane, Dean Rasheed)

This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.

- Fix executor crash due to double free in some GROUPING SET usages (Peter Geoghegan)

- Avoid crash if a table rewrite event trigger is added concurrently with a command that could call such a trigger (Álvaro Herrera, Andrew Gierth, Tom Lane)

- Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)

- Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)

- Fix possible leak or double free of visibility map buffer pins (Amit Kapila)

- Avoid spuriously marking pages as all-visible (Dan Wood, Pavan Deolasee, Álvaro Herrera)
This could happen if some tuples were locked (but not deleted). While queries would still function correctly, vacuum would normally ignore such pages, with the long-term effect that the tuples were never frozen. In recent releases this would eventually result in errors such as “found multixact nnnnn from before relminmxid nnnnn”.  

- Fix overly strict sanity check in heap_prepare_freeze_tuple (Álvaro Herrera)  
  This could result in incorrect “cannot freeze committed xmax” failures in databases that have been pg_upgrade'd from 9.2 or earlier.  
- Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)  
- Reduce locking during autovacuum worker scheduling (Jeff Janes)  
  The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.  
- Ensure client hostname is copied while copying pg_stat_activity data to local memory (Edmund Horner)  
  Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the client hostname column to change unexpectedly if any existing session disconnected.  
- Fix incorrect processing of multiple compound affixes in ispell dictionaries (Arthur Zakirov)  
- Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on text columns (Tom Lane)  
  Such searches would return the wrong set of rows in most non-C locales.  
- Prevent query-lifespan memory leakage with SP-GiST operator classes that use traversal values (Anton Dignös)  
- Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)  
  Previously, the tuple count was reported to be the same as that of the underlying table, which is wrong if the index is partial.  
- Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)  
  Previously it reported the estimated number of heap tuples, which might be inaccurate, and is certainly wrong if the index is partial.  
- Fix a corner case where a streaming standby gets stuck at a WAL continuation record (Kyotaro Horiguchi)  
- In logical decoding, avoid possible double processing of WAL data when a walsender restarts (Craig Ringer)  
- Allow scalarltsel and scalargtsel to be used on non-core datatypes (Tomas Vondra)  
- Reduce libpq's memory consumption when a server error is reported after a large amount of query output has been collected (Tom Lane)  
  Discard the previous output before, not after, processing the error message. On some platforms, notably Linux, this can make a difference in the application's subsequent memory footprint.  
- Fix double-free crashes in ecpg (Patrick Krecker, Jeevan Ladhe)  
- Fix ecpg to handle long long int variables correctly in MSVC builds (Michael Meskes, Andrew Gierth)
• Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)

The `local_preload_libraries`, `session_preload_libraries`, `shared_preload_libraries`, and `temp_tablespaces` variables were not correctly quoted in `pg_dump` output. This would cause problems if settings for these variables appeared in `CREATE FUNCTION ... SET` or `ALTER DATABASE/ROLE ... SET` clauses.

• Fix `pg_recvlogical` to not fail against pre-v10 PostgreSQL servers (Michael Paquier)

A previous fix caused `pg_recvlogical` to issue a command regardless of server version, but it should only be issued to v10 and later servers.

• Ensure that `pg_rewind` deletes files on the target server if they are deleted from the source server during the run (Takayuki Tsunakawa)

Failure to do this could result in data inconsistency on the target, particularly if the file in question is a WAL segment.

• Fix `pg_rewind` to handle tables in non-default tablespaces correctly (Takayuki Tsunakawa)

• Fix overflow handling in PL/pgSQL integer `FOR` loops (Tom Lane)

The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

• Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisentraut)

• Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)

• Fix errors in initial build of `contrib/bloom` indexes (Tomas Vondra, Tom Lane)

Fix possible omission of the table's last tuple from the index. Count the number of index tuples correctly, in case it is a partial index.

• Rename internal `b64_encode` and `b64_decode` functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)

• Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)

This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.

• Update time zone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

E.13. Release 9.6.8

Release date: 2018-03-01

This release contains a variety of fixes from 9.6.7. For information about new features in the 9.6 major release, see Section E.21.

E.13.1. Migration to Version 9.6.8

A dump/restore is not required for those running 9.6.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that
you read the documentation changes described in the first changelog entry below, and take suitable steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 9.6.7, see Section E.14.

E.13.2. Changes

- Document how to configure installations and applications to guard against search-path-dependent trojan-horse attacks from other users (Noah Misch)

Using a `search_path` setting that includes any schemas writable by a hostile user enables that user to capture control of queries and then run arbitrary SQL code with the permissions of the attacked user. While it is possible to write queries that are proof against such hijacking, it is notionally tedious, and it's very easy to overlook holes. Therefore, we now recommend configurations in which no untrusted schemas appear in one's search path. Relevant documentation appears in Section 5.8.6 (for database administrators and users), Section 34.1 (for application authors), Section 38.16.1 (for extension authors), and `CREATE FUNCTION` (for authors of `SECURITY DEFINER` functions). (CVE-2018-1058)

- Avoid use of insecure `search_path` settings in `pg_dump` and other client programs (Noah Misch, Tom Lane)

`pg_dump`, `pg_upgrade`, `vacuumdb` and other PostgreSQL-provided applications were themselves vulnerable to the type of hijacking described in the previous changelog entry; since these applications are commonly run by superusers, they present particularly attractive targets. To make them secure whether or not the installation as a whole has been secured, modify them to include only the `pg_catalog` schema in their `search_path` settings. Autovacuum worker processes now do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example, user-provided functions in index expressions — the tighter `search_path` may result in errors, which will need to be corrected by adjusting those user-provided functions to not assume anything about what search path they are invoked under. That has always been good practice, but now it will be necessary for correct behavior. (CVE-2018-1058)

- Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom Lane)

If a CTE (WITH clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

- Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do not match mergeclauses” planner errors in corner cases.

- Repair `pg_upgrade`'s failure to preserve `relfrozenxid` for materialized views (Tom Lane, Andres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The problem would be more likely to occur in seldom-refreshed materialized views, or ones that were maintained only with `REFRESH MATERIALIZED VIEW CONCURRENTLY`.

If such corruption is observed, it can be repaired by refreshing the materialized view (without `CONCURRENTLY`).
• Fix incorrect reporting of PL/Python function names in error CONTEXT stacks (Tom Lane)

An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function's name twice, rather than the expected results. Also, an error in a nested PL/Python DO block could result in a null pointer dereference crash on some platforms.

• Allow contrib/auto_explain's log_min_duration setting to range up to INT_MAX, or about 24 days instead of 35 minutes (Tom Lane)

• Mark assorted GUC variables as PGDLLIMPORT, to ease porting extension modules to Windows (Metin Doslu)


Release date: 2018-02-08

This release contains a variety of fixes from 9.6.6. For information about new features in the 9.6 major release, see Section E.21.

E.14.1. Migration to Version 9.6.7

A dump/restore is not required for those running 9.6.X.

However, if you use contrib/cube's -> operator, see the entry below about that.

Also, if you are upgrading from a version earlier than 9.6.6, see Section E.15.

E.14.2. Changes

• Ensure that all temporary files made by pg_upgrade are non-world-readable (Tom Lane, Noah Misch)

pg_upgrade normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing pg_dumpall -q output would be group- or world-readable, or even writable, if the user's umask setting allows. In typical usage on multi-user machines, the umask and/or the working directory's permissions would be tight enough to prevent problems; but there may be people using pg_upgrade in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

• Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

In some cases VACUUM would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

• Ensure that vacuum will always clean up the pending-insertions list of a GIN index (Masahiko Sawada)

This is necessary to ensure that dead index entries get removed. The old code got it backwards, allowing vacuum to skip the cleanup if some other process were running cleanup concurrently, thus risking invalid entries being left behind in the index.

• Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)

These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.

• Fix incorrect query results from cases involving flattening of subqueries whose outputs are used in GROUPING SETS (Heikki Linnakangas)
• Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by ALTER TABLE NO INHERIT (Tom Lane)

• Fix spurious deadlock failures when multiple sessions are running CREATE INDEX CONCURRENTLY (Jeff Janes)

• Fix failures when an inheritance tree contains foreign child tables (Etsuro Fujita)

A mix of regular and foreign tables in an inheritance tree resulted in creation of incorrect plans for UPDATE and DELETE queries. This led to visible failures in some cases, notably when there are row-level triggers on a foreign child table.

• Repair failure with correlated sub-SELECT inside VALUES inside a LATERAL subquery (Tom Lane)

• Fix “could not devise a query plan for the given query” planner failure for some cases involving nested UNION ALL inside a lateral subquery (Tom Lane)

• Fix logical decoding to correctly clean up disk files for crashed transactions (Atsushi Torikoshi)

Logical decoding may spill WAL records to disk for transactions generating many WAL records. Normally these files are cleaned up after the transaction's commit or abort record arrives; but if no such record is ever seen, the removal code misbehaved.

• Fix walsender timeout failure and failure to respond to interrupts when processing a large transaction (Petr Jelinek)

• Fix has_sequence_privilege() to support WITH GRANT OPTION tests, as other privilege-testing functions do (Joe Conway)

• In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)

We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don't promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only xpath() and related functions; other XML code paths already acted this way.

• Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)

Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.

• Cope with failure to start a parallel worker process (Amit Kapila, Robert Haas)

Parallel query previously tended to hang indefinitely if a worker could not be started, as the result of fork() failure or other low-probability problems.

• Fix collection of EXPLAIN statistics from parallel workers (Amit Kapila, Thomas Munro)

• Avoid unsafe alignment assumptions when working with __int128 (Tom Lane)

Typically, compilers assume that __int128 variables are aligned on 16-byte boundaries, but our memory allocation infrastructure isn't prepared to guarantee that, and increasing the setting of MAXALIGN seems infeasible for multiple reasons. Adjust the code to allow use of __int128 only when we can tell the compiler to assume lesser alignment. The only known symptom of this problem so far is crashes in some parallel aggregation queries.
Release Notes

- Prevent stack-overflow crashes when planning extremely deeply nested set operations (UNION/INTERSECT/EXCEPT) (Tom Lane)

- Fix null-pointer crashes for some types of LDAP URLs appearing in pg_hba.conf (Thomas Munro)

- Fix sample INSTR() functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)

  These functions are stated to be Oracle® compatible, but they weren't exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.

  The sample code has been adjusted to match Oracle's behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

- Fix pg_dump to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

  The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with DATABASE so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with LARGE OBJECT. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.

  Note that this change has user-visible results in the output of pg_restore --list.

- Rename pg_rewind's copy_file_range function to avoid conflict with new Linux system call of that name (Andres Freund)

  This change prevents build failures with newer glibc versions.

- In ecpg, detect indicator arrays that do not have the correct length and report an error (David Rader)

- Change the behavior of contrib/cube's cube ~> int operator to make it compatible with KNN search (Alexander Korotkov)

  The meaning of the second argument (the dimension selector) has been changed to make it predictable which value is selected even when dealing with cubes of varying dimensionalities.

  This is an incompatible change, but since the point of the operator was to be used in KNN searches, it seems rather useless as-is. After installing this update, any expression indexes or materialized views using this operator will need to be reindexed/refreshed.

- Avoid triggering a libc assertion in contrib/hstore, due to use of memcpy() with equal source and destination pointers (Tomas Vondra)

- Fix incorrect display of tuples' nullBitmaps in contrib/pageinspect (Maksim Milyutin)

- In contrib/postgres_fdw, avoid “outer pathkeys do not match mergeclauses” planner error when constructing a plan involving a remote join (Robert Haas)

- Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

  The scripts in contrib/start-scripts/osx use infrastructure that's been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory contrib/start-scripts/macos containing scripts that use the newer launchd infrastructure.
• Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dun-stand)

• Support linking to MinGW-built versions of libperl (Noah Misch)

  This allows building PL/Perl with some common Perl distributions for Windows.

• Fix MSVC build to test whether 32-bit libperl needs -D_USE_32BIT_TIME_T (Noah Misch)

  Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.

• On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)

  This may allow collection of a core dump for some early-startup failures that did not produce a dump before.

• On Windows, avoid encoding-conversion-related crashes when emitting messages very early in postmaster startup (Takayuki Tsunakawa)

• Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)

• Add support for spinlocks on Motorola 88K (David Carlier)

• Update time zone data files to tzdata release 2018c for DST law changes in Brazil, Sao Tome and Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The US/Pacific-New zone has been removed (it was only an alias for America/Los_Angeles anyway).

E.15. Release 9.6.6

Release date: 2017-11-09

This release contains a variety of fixes from 9.6.5. For information about new features in the 9.6 major release, see Section E.21.

E.15.1. Migration to Version 9.6.6

A dump/restore is not required for those running 9.6.X.

However, if you use BRIN indexes, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.6.4, see Section E.17.

E.15.2. Changes

• Ensure that INSERT ... ON CONFLICT DO UPDATE checks table permissions and RLS policies in all cases (Dean Rasheed)

  The update path of INSERT ... ON CONFLICT DO UPDATE requires SELECT permission on the columns of the arbiter index, but it failed to check for that in the case of an arbiter specified by constraint name. In addition, for a table with row level security enabled, it failed to check updated rows against the table's SELECT policies (regardless of how the arbiter index was specified). (CVE-2017-15099)

• Fix crash due to rowtype mismatch in json{b}_populate_recordset() (Michael Paquier, Tom Lane)

  These functions used the result rowtype specified in the FROM ... AS clause without checking that it matched the actual rowtype of the supplied tuple value. If it didn't, that would usually result in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)
• Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)

Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making $PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

• Fix BRIN index summarization to handle concurrent table extension correctly (Álvaro Herrera)

Previously, a race condition allowed some table rows to be omitted from the index. It may be necessary to reindex existing BRIN indexes to recover from past occurrences of this problem.

• Fix possible failures during concurrent updates of a BRIN index (Tom Lane)

These race conditions could result in errors like “invalid index offnum” or “inconsistent range map”.

• Fix crash when logical decoding is invoked from a SPI-using function, in particular any function written in a PL language (Tom Lane)

• Fix incorrect query results when multiple GROUPING SETS columns contain the same simple variable (Tom Lane)

• Fix incorrect parallelization decisions for nested queries (Amit Kapila, Kuntal Ghosh)

• Fix parallel query handling to not fail when a recently-used role is dropped (Amit Kapila)

• Fix json_build_array(), json_build_object(), and their jsonb equivalents to handle explicit VARIADIC arguments correctly (Michael Paquier)

• Properly reject attempts to convert infinite float values to type numeric (Tom Lane, KaiGai Kohei)

Previously the behavior was platform-dependent.

• Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)

• Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes (Tom Lane)

Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

• Correctly detect hashability of range data types (Tom Lane)

The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range's subtype has hash support. This does not affect any of the built-in range types, since they're all hashable anyway.

• Correctly ignore RelabelType expression nodes when determining relation distinctness (David Rowley)

This allows the intended optimization to occur when a subquery has a result column of type varchar.

• Prevent sharing transition states between ordered-set aggregates (David Rowley)

This causes a crash with the built-in ordered-set aggregates, and probably with user-written ones as well. v11 and later will include provisions for dealing with such cases safely, but in released branches, just disable the optimization.
• Prevent `idle_in_transaction_session_timeout` from being ignored when a statement_timeout occurred earlier (Lukas Fittl)

• Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)

  If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.

• Avoid SIGBUS crash on Linux when a DSM memory request exceeds the space available in tmpfs (Thomas Munro)

• Reduce the frequency of data flush requests during bulk file copies to avoid performance problems on macOS, particularly with its new APFS file system (Tom Lane)

• Prevent low-probability crash in processing of nested trigger firings (Tom Lane)

• Allow `COPY`'s FREEZE option to work when the transaction isolation level is REPEATABLE READ or higher (Noah Misch)

  This case was unintentionally broken by a previous bug fix.

• Correctly restore the umask setting when file creation fails in `COPY` or `lo_export()` (Peter Eisentraut)

• Give a better error message for duplicate column names in `ANALYZE` (Nathan Bossart)

• Add missing cases in `GetCommandLogLevel()`, preventing errors when certain SQL commands are used while `log_statement` is set to `ddl` (Michael Paquier)

• Fix mis-parsing of the last line in a non-newline-terminated `pg_hba.conf` file (Tom Lane)

• Fix `AggGetAggref()` to return the correct `Aggref` nodes to aggregate final functions whose transition calculations have been merged (Tom Lane)

• Fix `pg_dump` to ensure that it emits `GRANT` commands in a valid order (Stephen Frost)

• Fix `pg_basebackup`'s matching of tablespace paths to canonicalize both paths before comparing (Michael Paquier)

  This is particularly helpful on Windows.

• Fix `libpq` to not require user's home directory to exist (Tom Lane)

  In v10, failure to find the home directory while trying to read `~/.pgpass` was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading `~/.pg_service.conf`, though this was less obvious since that file is not sought unless a service name is specified.

• Fix `libpq` to guard against integer overflow in the row count of a `PGresult` (Michael Paquier)

• Fix `ecpg`'s handling of out-of-scope cursor declarations with pointer or array variables (Michael Meskes)

• In `ecpglib`, correctly handle backslashes in string literals depending on whether `standard_conforming_strings` is set (Tsunakawa Takayuki)

• Make `ecpglib`'s Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengqi, Michael Meskes)

• Fix `ecpg`'s regression tests to work reliably on Windows (Christian Ullrich, Michael Meskes)

• Fix missing temp-install prerequisites for `check`-like Make targets (Noah Misch)
Some non-default test procedures that are meant to work like `make check` failed to ensure that the temporary installation was up to date.

- Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

  This fixes various issues: the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no `posixrules` file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.

- Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

## E.16. Release 9.6.5

**Release date:** 2017-08-31

This release contains a small number of fixes from 9.6.4. For information about new features in the 9.6 major release, see Section E.21.

### E.16.1. Migration to Version 9.6.5

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.4, see Section E.17.

### E.16.2. Changes

- Show foreign tables in `information_schema.table_privileges` view (Peter Eisentraut)

  All other relevant `information_schema` views include foreign tables, but this one ignored them.

  Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:

```sql
SET search_path TO information_schema;
CREATE OR REPLACE VIEW table_privileges AS
  SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
         CAST(grantee.rolname AS sql_identifier) AS grantee,
         CAST(current_database() AS sql_identifier) AS table_catalog,
         CAST(nc.nspname AS sql_identifier) AS table_schema,
         CAST(c.relname AS sql_identifier) AS table_name,
         CAST(c.prtype AS character_data) AS privilege_type,
         CAST(
           CASE WHEN -- object owner always has grant options
                   pg_has_role(grantee.oid, c.relowner, 'USAGE')
           OR c.grantable
           THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
         CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE
                     'NO' END AS yes_or_no) AS with_hierarchy
  FROM (2442
```
SELECT oid, relname, relnamespace, relkind, relowner, (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class
) AS c (oid, relname, relnamespace, relkind, relowner, grantor, grantee, prtype, grantable),
pg_namespace nc,
pg_authid u_grantor,
(
SELECT oid, rolname FROM pg_authid
UNION ALL
SELECT 0::oid, 'PUBLIC'
) AS grantee (oid, rolname)
WHERE c.relnamespace = nc.oid
AND c.relkind IN ('r', 'v', 'f')
AND c.grantee = grantee.oid
AND c.grantor = u_grantor.oid
AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE', 'DELETE', 'TRUNCATE', 'REFERENCES', 'TRIGGER')
AND (pg_has_role(u_grantor.oid, 'USAGE')
OR pg_has_role(grantee.oid, 'USAGE')
OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

• Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)

This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about “cannot drop active portal”.

• Remove assertion that could trigger during a fatal exit (Tom Lane)

• Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant values that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.

• Prevent crash when passing fixed-length pass-by-reference data types to parallel worker processes (Tom Lane)

• Fix crash in pg_restore when using parallel mode and using a list file to select a subset of items to restore (Fabrizio de Royes Mello)

• Change ecpg's parser to allow RETURNING clauses without attached C variables (Michael Meskes)

This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

• Change ecpg's parser to recognize backslash continuation of C preprocessor command lines (Michael Meskes)

• Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that
_USE_32BIT_TIME_T is no longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

- Fix make check to behave correctly when invoked via a non-GNU make program (Thomas Munro)

## E.17. Release 9.6.4

**Release date:** 2017-08-10

This release contains a variety of fixes from 9.6.3. For information about new features in the 9.6 major release, see Section E.21.

### E.17.1. Migration to Version 9.6.4

A dump/restore is not required for those running 9.6.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.6.3, see Section E.18.

### E.17.2. Changes

- Further restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Noah Misch)

The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have USAGE permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since information_schema.user_mapping_options does not show the options in such cases, pg_user_mappings should not either. (CVE-2017-7547)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, you will need to do the following:

1. Restart the postmaster after adding allow_system_table_mods = true to postgresql.conf. (In versions supporting ALTER SYSTEM, you can use that to make the configuration change, but you’ll still need a restart.)

2. In each database of the cluster, run the following commands as superuser:

   ```sql
   SET search_path = pg_catalog;
   CREATE OR REPLACE VIEW pg_user_mappings AS
   SELECT
     U.oid       AS umid,
     S.oid       AS srvid,
     S.srvname   AS srvname,
     U.umuser    AS umuser,
     CASE WHEN U.umuser = 0 THEN 'public'
     ELSE A.rolname END AS usename,
     CASE WHEN (U.umuser <> 0 AND A.rolname = current_user AND (pg_has_role(S.srvowner, 'USAGE') OR has_server_privilege(S.oid, 'USAGE')))
   ```
SELECT (CASE WHEN U.umuser = 0 AND pg_has_role(S.srvowner, 'USAGE')
               OR (SELECT rolsuper FROM pg_authid WHERE rolname = current_user)
               THEN U.umoptions
               ELSE NULL END AS umoptions
FROM pg_user_mapping U
LEFT JOIN pg_authid A ON (A.oid = U.umuser) JOIN pg_foreign_server S ON (U.umserver = S.oid);

3. Do not forget to include the template0 and template1 databases, or the vulnerability will still exist in databases you create later. To fix template0, you'll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;

and then after fixing template0, undo that with

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;

In prior versions, instead use

UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';

4. Finally, remove the allow_system_table_mods configuration setting, and again restart the postmaster.

- Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)

libpq ignores empty password specifications, and does not transmit them to the server. So, if a user's password has been set to the empty string, it's impossible to log in with that password via psql or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, md5, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

- Make lo_put() check for UPDATE privilege on the target large object (Tom Lane, Michael Paquier)

lo_put() should surely require the same permissions as lowrite(), but the check was missing, allowing any user to change the data in a large object. (CVE-2017-7548)

- Correct the documentation about the process for upgrading standby servers with pg_upgrade (Bruce Momjian)

The previous documentation instructed users to start/stop the primary server after running pg_upgrade but before syncing the standby servers. This sequence is unsafe.

- Fix concurrent locking of tuple update chains (Álvaro Herrera)

If several sessions concurrently lock a tuple update chain with nonconflicting lock modes using an old snapshot, and they all succeed, it was possible for some of them to nonetheless fail (and conclude there is no live tuple version) due to a race condition. This had consequences such as foreign-key checks failing to see a tuple that definitely exists but is being updated concurrently.

2445
• Fix potential data corruption when freezing a tuple whose XMAX is a multixact with exactly one still-interesting member (Teodor Sigaev)

• Avoid integer overflow and ensuing crash when sorting more than one billion tuples in-memory (Sergey Koposov)

• On Windows, retry process creation if we fail to reserve the address range for our shared memory in the new process (Tom Lane, Amit Kapila)

  This is expected to fix infrequent child-process-launch failures that are probably due to interference from antivirus products.

• Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas Munro, Tom Lane)

• Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael Paquier)

• Prevent sending SSL session tickets to clients (Tom Lane)

  This fix prevents reconnection failures with ticket-aware client-side SSL code.

• Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)

• Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and immediate restart (Tom Lane)

  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively ignored.

• Ensure that the statistics collector's receive buffer size is at least 100KB (Tom Lane)

  This reduces the risk of dropped statistics data on older platforms whose default receive buffer size is less than that.

• Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes an XLOG_SWITCH WAL record (Andres Freund)

• Fix walsender to exit promptly when client requests shutdown (Tom Lane)

• Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)

• Prevent walsender-triggered panics during shutdown checkpoints (Andres Freund, Michael Paquier)

• Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom Lane)

• Fix leakage of small subtransactions spilled to disk during logical decoding (Andres Freund)

  This resulted in temporary files consuming excessive disk space.

• Reduce the work needed to build snapshots during creation of logical-decoding slots (Andres Freund, Petr Jelinek)

  The previous algorithm was infeasibly expensive on a server with a lot of open transactions.

• Fix race condition that could indefinitely delay creation of logical-decoding slots (Andres Freund, Petr Jelinek)

• Reduce overhead in processing syscache invalidation events (Tom Lane)

  This is particularly helpful for logical decoding, which triggers frequent cache invalidation.
Release Notes

- Remove incorrect heuristic used in some cases to estimate join selectivity based on the presence of foreign-key constraints (David Rowley)

In some cases where a multi-column foreign key constraint existed but did not exactly match a query's join structure, the planner used an estimation heuristic that turns out not to work well at all. Revert such cases to the way they were estimated before 9.6.

- Fix cases where an INSERT or UPDATE assigns to more than one element of a column that is of domain-over-array type (Tom Lane)

- Allow window functions to be used in sub-SELECTs that are within the arguments of an aggregate function (Tom Lane)

- Ensure that a view's CHECK OPTIONS clause is enforced properly when the underlying table is a foreign table (Etsuro Fujita)

Previously, the update might get pushed entirely to the foreign server, but the need to verify the view conditions was missed if so.

- Move autogenerated array types out of the way during ALTER ... RENAME (Vik Fearing)

Previously, we would rename a conflicting autogenerated array type out of the way during CREATE; this fix extends that behavior to renaming operations.

- Fix dangling pointer in ALTER TABLE when there is a comment on a constraint belonging to the table (David Rowley)

Re-applying the comment to the reconstructed constraint could fail with a weird error message, or even crash.

- Ensure that ALTER USER ... SET accepts all the syntax variants that ALTER ROLE ... SET does (Peter Eisentraut)

- Allow a foreign table's CHECK constraints to be initially NOT VALID (Amit Langote)

CREATE TABLE silently drops NOT VALID specifiers for CHECK constraints, reasoning that the table must be empty so the constraint can be validated immediately. But this is wrong for CREATE FOREIGN TABLE, where there's no reason to suppose that the underlying table is empty, and even if it is it's no business of ours to decide that the constraint can be treated as valid going forward. Skip this "optimization" for foreign tables.

- Properly update dependency info when changing a datatype I/O function's argument or return type from opaque to the correct type (Heikki Linnakangas)

CREATE TYPE updates I/O functions declared in this long-obsolete style, but it forgot to record a dependency on the type, allowing a subsequent DROP TYPE to leave broken function definitions behind.

- Allow parallelism in the query plan when COPY copies from a query's result (Andres Freund)

- Reduce memory usage when ANALYZE processes a tsvector column (Heikki Linnakangas)

- Fix unnecessary precision loss and sloppy rounding when multiplying or dividing money values by integers or floats (Tom Lane)

- Tighten checks for whitespace in functions that parse identifiers, such as regprocedurein() (Tom Lane)

Depending on the prevailing locale, these functions could misinterpret fragments of multibyte characters as whitespace.

- Use relevant #define symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)
This avoids portability problems, typically manifesting as a “handshake” mismatch during library load, when working with recent Perl versions.

- In libpq, reset GSS/SASL and SSPI authentication state properly after a failed connection attempt (Michael Paquier)

Failure to do this meant that when falling back from SSL to non-SSL connections, a GSS/SASL failure in the SSL attempt would always cause the non-SSL attempt to fail. SSPI did not fail, but it leaked memory.

- In psql, fix failure when COPY FROM STDIN is ended with a keyboard EOF signal and then another COPY FROM STDIN is attempted (Thomas Munro)

This misbehavior was observed on BSD-derived platforms (including macOS), but not on most others.

- Fix pg_dump and pg_restore to emit REFRESH MATERIALIZED VIEW commands last (Tom Lane)

This prevents errors during dump/restore when a materialized view refers to tables owned by a different user.

- Improve pg_dump/pg_restore's reporting of error conditions originating in zlib (Vladimir Kunschikov, Álvaro Herrera)

- Fix pg_dump with the --clean option to drop event triggers as expected (Tom Lane)

It also now correctly assigns ownership of event triggers; before, they were restored as being owned by the superuser running the restore script.

- Fix pg_dump with the --clean option to not fail when the public schema doesn't exist (Stephen Frost)

- Fix pg_dump to not emit invalid SQL for an empty operator class (Daniel Gustafsson)

- Fix pg_dump output to stdout on Windows (Kuntal Ghosh)

A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_get_ruledef() to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)

In some corner cases, pg_dump relies on pg_get_ruledef() to dump views, so that this error could result in dump/reload failures.

- Fix dumping of outer joins with empty constraints, such as the result of a NATURAL LEFT JOIN with no common columns (Tom Lane)

- Fix dumping of function expressions in the FROM clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)

- Fix pg_basebackup output to stdout on Windows (Haribabu Kommi)

A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_REWIND to correctly handle files exceeding 2GB (Kuntal Ghosh, Michael Paquier)

Ordinarily such files won't appear in PostgreSQL data directories, but they could be present in some cases.
• Fix pg_upgrade to ensure that the ending WAL record does not have wal_level = minimum (Bruce Momjian)

This condition could prevent upgraded standby servers from reconnecting.

• Fix pg_xlogdump’s computation of WAL record length (Andres Freund)

• In postgres_fdw, re-establish connections to remote servers after ALTER SERVER or ALTER USER MAPPING commands (Kyotaro Horiguchi)

This ensures that option changes affecting connection parameters will be applied promptly.

• In postgres_fdw, allow cancellation of remote transaction control commands (Robert Haas, Rafia Sabih)

This change allows us to quickly escape a wait for an unresponsive remote server in many more cases than previously.

• Increase MAX SYSCACHE CALLBACKS to provide more room for extensions (Tom Lane)

• Always use -fPIC, not -fpic, when building shared libraries with gcc (Tom Lane)

This supports larger extension libraries on platforms where it makes a difference.

• In MSVC builds, handle the case where the OpenSSL library is not within a VC subdirectory (Andrew Dunstan)

• In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)

This fixes a former need to move things around in standard Windows installations of libxml2.

• In MSVC builds, recognize a Tcl library that is named tcl86.lib (Noah Misch)

• In MSVC builds, honor PROVE_FLAGS settings on vcregress.pl's command line (Andrew Dunstan)

E.18. Release 9.6.3

Release date: 2017-05-11

This release contains a variety of fixes from 9.6.2. For information about new features in the 9.6 major release, see Section E.21.

E.18.1. Migration to Version 9.6.3

A dump/restore is not required for those running 9.6.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are using third-party replication tools that depend on “logical decoding”, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.6.2, see Section E.19.

E.18.2. Changes

• Restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)

The previous coding allowed the owner of a foreign server object, or anyone he has granted server USAGE permission to, to see the options for all user mappings associated with that server. This
might well include passwords for other users. Adjust the view definition to match the behavior of `information_schema.user_mapping_options`, namely that these options are visible to the user being mapped, or if the mapping is for `PUBLIC` and the current user is the server owner, or if the current user is a superuser. (CVE-2017-7486)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, follow the corrected procedure shown in the changelog entry for CVE-2017-7547, in Section E.17.

- Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

Some selectivity estimation functions in the planner will apply user-defined operators to values obtained from `pg_statistic`, such as most common values and histogram entries. This occurs before table permissions are checked, so a nefarious user could exploit the behavior to obtain these values for table columns he does not have permission to read. To fix, fall back to a default estimate if the operator's implementation function is not certified leak-proof and the calling user does not have permission to read the table column whose statistics are needed. At least one of these criteria is satisfied in most cases in practice. (CVE-2017-7484)

- Restore libpq's recognition of the `PGREQUIRESSL` environment variable (Daniel Gustafsson)

Processing of this environment variable was unintentionally dropped in PostgreSQL 9.3, but its documentation remained. This creates a security hazard, since users might be relying on the environment variable to force SSL-encrypted connections, but that would no longer be guaranteed. Restore handling of the variable, but give it lower priority than `PGSSLMODE`, to avoid breaking configurations that work correctly with post-9.3 code. (CVE-2017-7485)

- Fix possibly-invalid initial snapshot during logical decoding (Petr Jelinek, Andres Freund)

The initial snapshot created for a logical decoding replication slot was potentially incorrect. This could cause third-party tools that use logical decoding to copy incomplete/inconsistent initial data. This was more likely to happen if the source server was busy at the time of slot creation, or if another logical slot already existed.

If you are using a replication tool that depends on logical decoding, and it should have copied a nonempty data set at the start of replication, it is advisable to recreate the replica after installing this update, or to verify its contents against the source server.

- Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)

This could result in an unlogged index being set to an invalid state after a crash and restart. Such a problem would persist until the index was dropped and rebuilt.

- Fix incorrect reconstruction of `pg_subtrans` entries when a standby server replays a prepared but uncommitted two-phase transaction (Tom Lane)

In most cases this turned out to have no visible ill effects, but in corner cases it could result in circular references in `pg_subtrans`, potentially causing infinite loops in queries that examine rows modified by the two-phase transaction.

- Avoid possible crash in walsender due to failure to initialize a string buffer (Stas Kelvich, Fujii Masao)

- Fix possible crash when rescanning a nearest-neighbor index-only scan on a GiST index (Tom Lane)

- Prevent delays in postmaster's launching of multiple parallel worker processes (Tom Lane)

There could be a significant delay (up to tens of seconds) before satisfying a query’s request for more than one worker process, or when multiple queries requested workers simultaneously. On most platforms this required unlucky timing, but on some it was the typical case.
• Fix postmaster’s handling of fork() failure for a background worker process (Tom Lane)

Previously, the postmaster updated portions of its state as though the process had been launched successfully, resulting in subsequent confusion.

• Fix possible “no relation entry for relid 0” error when planning nested set operations (Tom Lane)

• Fix assorted minor issues in planning of parallel queries (Robert Haas)

• Avoid applying “physical targetlist” optimization to custom scans (Dmitry Ivanov, Tom Lane)

This optimization supposed that retrieving all columns of a tuple is inexpensive, which is true for ordinary Postgres tuples; but it might not be the case for a custom scan provider.

• Use the correct sub-expression when applying a FOR ALL row-level-security policy (Stephen Frost)

In some cases the WITH CHECK restriction would be applied when the USING restriction is more appropriate.

• Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL (Julien Rouhaud, Tom Lane)

Due to lack of a cache flush step between commands in an extension script file, non-utility queries might not see the effects of an immediately preceding catalog change, such as ALTER TABLE ... RENAME.

• Skip tablespace privilege checks when ALTER TABLE ... ALTER COLUMN TYPE rebuilds an existing index (Noah Misch)

The command failed if the calling user did not currently have CREATE privilege for the tablespace containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be rebuilt where it is.

• Fix ALTER TABLE ... VALIDATE CONSTRAINT to not recurse to child tables when the constraint is marked NO INHERIT (Amit Langote)

This fix prevents unwanted “constraint does not exist” failures when no matching constraint is present in the child tables.

• Avoid dangling pointer in COPY ... TO when row-level security is active for the source table (Tom Lane)

Usually this had no ill effects, but sometimes it would cause unexpected errors or crashes.

• Avoid accessing an already-closed relcache entry in CLUSTER and VACUUM FULL (Tom Lane)

With some bad luck, this could lead to indexes on the target relation getting rebuilt with the wrong persistence setting.

• Fix VACUUM to account properly for pages that could not be scanned due to conflicting page pins (Andrew Gierth)

This tended to lead to underestimation of the number of tuples in the table. In the worst case of a small heavily-contended table, VACUUM could incorrectly report that the table contained no tuples, leading to very bad planning choices.

• Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests (Tom Lane, Thomas Munro)

• Fix incorrect support for certain box operators in SP-GiST (Nikita Glukhov)

SP-GiST index scans using the operators &< &> &< | and | &> would yield incorrect answers.
• Fix integer-overflow problems in `interval` comparison (Kyotaro Horiguchi, Tom Lane)

The comparison operators for type `interval` could yield wrong answers for intervals larger than about 296000 years. Indexes on columns containing such large values should be reindexed, since they may be corrupt.

• Fix `cursor_to_xml()` to produce valid output with `tableforest = false` (Thomas Munro, Peter Eisentraut)

Previously it failed to produce a wrapping `<table>` element.

• Fix roundoff problems in `float8_timestamptz()` and `make_interval()` (Tom Lane)

These functions truncated, rather than rounded, when converting a floating-point value to integer microseconds; that could cause unexpectedly off-by-one results.

• Fix `pg_get_object_address()` to handle members of operator families correctly (Álvaro Herrera)

• Fix cancelling of `pg_stop_backup()` when attempting to stop a non-exclusive backup (Michael Paquier, David Steele)

If `pg_stop_backup()` was cancelled while waiting for a non-exclusive backup to end, related state was left inconsistent; a new exclusive backup could not be started, and there were other minor problems.

• Improve performance of `pg_timezone_names` view (Tom Lane, David Rowley)

• Reduce memory management overhead for contexts containing many large blocks (Tom Lane)

• Fix sloppy handling of corner-case errors from `lseek()` and `close()` (Tom Lane)

Neither of these system calls are likely to fail in typical situations, but if they did, `fd.c` could get quite confused.

• Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)

This could result in attempting to write to the event log when that isn't accessible, so that no logging happens at all.

• Fix ecpg to support `COMMIT PREPARED` and `ROLLBACK PREPARED` (Masahiko Sawada)

• Fix a double-free error when processing dollar-quoted string literals in ecpg (Michael Meskes)

• Fix pgbench to handle the combination of `--connect` and `--rate` options correctly (Fabien Coelho)

• Fix pgbench to honor the long-form option spelling `--builtin`, as per its documentation (Tom Lane)

• Fix `pg_dump/pg_restore` to correctly handle privileges for the `public` schema when using `--clean` option (Stephen Frost)

Other schemas start out with no privileges granted, but `public` does not; this requires special-case treatment when it is dropped and restored due to the `--clean` option.

• In `pg_dump`, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)

In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.
• Fix typo in pg_dump's query for initial privileges of a procedural language (Peter Eisentraut)

This resulted in pg_dump always believing that the language had no initial privileges. Since that's true for most procedural languages, ill effects from this bug are probably rare.

• Avoid emitting an invalid list file in pg_restore -l when SQL object names contain newlines (Tom Lane)

Replace newlines by spaces, which is sufficient to make the output valid for pg_restore -L's purposes.

• Fix pg_upgrade to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)

Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.

• Improve error handling in contrib/adminpack's pg_file_write() function (Noah Misch)

Notably, it failed to detect errors reported by fclose().

• In contrib/dblink, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

• Fix contrib/pg_trgm's extraction of trigrams from regular expressions (Tom Lane)

In some cases it would produce a broken data structure that could never match anything, leading to GIN or GiST indexscans that use a trigram index not finding any matches to the regular expression.

• In contrib/postgres_fdw, allow join conditions that contain shippable extension-provided functions to be pushed to the remote server (David Rowley, Ashutosh Bapat)

• Support Tcl 8.6 in MSVC builds (Álvaro Herrera)

• Sync our copy of the timezone library with IANA release tzcode2017b (Tom Lane)

This fixes a bug affecting some DST transitions in January 2038.

• Update time zone data files to tzdata release 2017b for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

• Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds (David Rowley)

The Microsoft MSVC build scripts neglected to install the posixrules file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the US/Eastern zone. If you don't want that, remove the posixrules file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.
E.19. Release 9.6.2

Release date: 2017-02-09

This release contains a variety of fixes from 9.6.1. For information about new features in the 9.6 major release, see Section E.21.

E.19.1. Migration to Version 9.6.2

A dump/restore is not required for those running 9.6.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.6.1, see Section E.20.

E.19.2. Changes

• Fix a race condition that could cause indexes built with CREATE INDEX CONCURRENTLY to be corrupt (Pavan Deolasee, Tom Lane)

If CREATE INDEX CONCURRENTLY was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the CREATE INDEX command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

• Ensure that the special snapshot used for catalog scans is not invalidated by premature data pruning (Tom Lane)

Backends failed to account for this snapshot when advertising their oldest xmin, potentially allowing concurrent vacuuming operations to remove data that was still needed. This led to transient failures along the lines of “cache lookup failed for relation 1255”.

• Fix incorrect WAL logging for BRIN indexes (Kuntal Ghosh)

The WAL record emitted for a BRIN “revmap” page when moving an index tuple to a different page was incorrect. Replay would make the related portion of the index useless, forcing it to be recomputed.

• Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)

Previously, this was skipped when wal_level = minimal, but actually it's necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

• If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)

• Ensure that hot standby feedback works correctly when it's enabled at standby server start (Ants Aasma, Craig Ringer)

• Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

• Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)

This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

• Disallow setting the num_sync field to zero in synchronous_standby_names (Fujii Masao)

The correct way to disable synchronous standby is to set the whole value to an empty string.
• Don’t count background worker processes against a user’s connection limit (David Rowley)

• Fix check for when an extension member object can be dropped (Tom Lane)

  Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

• Fix tracking of initial privileges for extension member objects so that it works correctly with ALTER EXTENSION ... ADD/DROP (Stephen Frost)

  An object’s current privileges at the time it is added to the extension will now be considered its default privileges; only later changes in its privileges will be dumped by subsequent pg_dump runs.

• Make sure ALTER TABLE preserves index tablespace assignments when rebuilding pg_dump runs (Tom Lane, Michael Paquier)

  Previously, non-default settings of default_tablespace could result in weird behavior.

• Fix incorrect updating of trigger function properties when changing a foreign-key constraint’s deferrability properties with ALTER TABLE ... ALTER CONSTRAINT (Tom Lane)

  This led to odd failures during subsequent exercise of the foreign key, as the triggers were fired at the wrong times.

• Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)

  This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.

• Fix ALTER TABLE ... SET DATA TYPE ... USING when child table has different column ordering than the parent (Álvaro Herrera)

  Failure to adjust the column numbering in the USING expression led to errors, typically “attribute N has wrong type”.

• Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via ALTER TABLE ... INHERIT (Amit Langote)

  The OID column should be treated the same as regular user columns in this case, but it wasn’t, leading to odd behavior in later inheritance changes.

• Ensure that CREATE TABLE ... LIKE ... WITH OIDS creates a table with OIDs, whether or not the LIKE-referenced table(s) have OIDs (Tom Lane)

• Fix CREATE OR REPLACE VIEW to update the view query before attempting to apply the new view options (Dean Rasheed)

  Previously the command would fail if the new options were inconsistent with the old view definition.

• Report correct object identity during ALTER TEXT SEARCH CONFIGURATION (Artur Zakirov)

  The wrong catalog OID was reported to extensions such as logical decoding.

• Fix commit timestamp mechanism to not fail when queried about the special XIDs Frozen-TransactionId and BootstrapTransactionId (Craig Ringer)

• Fix incorrect use of view reloptions as regular table reloptions (Tom Lane)

  The symptom was spurious “ON CONFLICT is not supported on table ... used as a catalog table” errors when the target of INSERT ... ON CONFLICT is a view with cascade option.

• Fix incorrect “target lists can have at most N entries” complaint when using ON CONFLICT with wide tables (Tom Lane)
• Fix spurious “query provides a value for a dropped column” errors during INSERT or UPDATE on a table with a dropped column (Tom Lane)

• Prevent multicolumn expansion of foo.* in an UPDATE source expression (Tom Lane)

  This led to “UPDATE target count mismatch --- internal error”. Now the syntax is understood as a whole-row variable, as it would be in other contexts.

• Ensure that column typmods are determined accurately for multi-row VALUES constructs (Tom Lane)

  This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a varchar value) but later values don't share the same limit.

• Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

  Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (U&'...') or Unicode identifier (U&"...").

• Fix execution of DISTINCT and ordered aggregates when multiple such aggregates are able to share the same transition state (Heikki Linnakangas)

• Fix implementation of phrase search operators in tsquery (Tom Lane)

  Remove incorrect, and inconsistently-applied, rewrite rules that tried to transform away AND/OR/NOT operators appearing below a PHRASE operator; instead upgrade the execution engine to handle such cases correctly. This fixes assorted strange behavior and possible crashes for text search queries containing such combinations. Also fix nested PHRASE operators to work sanely in combinations other than simple left-deep trees, correct the behavior when removing stopwords from a phrase search clause, and make sure that index searches behave consistently with simple sequential-scan application of such queries.

• Ensure that a purely negative text search query, such as !foo, matches empty tsvector values (Tom Dunstan)

  Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

• Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

• Fix performance problems in ts_rewrite() (Tom Lane)

• Fix ts_rewrite()'s handling of nested NOT operators (Tom Lane)

• Improve speed of user-defined aggregates that use array_append() as transition function (Tom Lane)

• Fix array_fill() to handle empty arrays properly (Tom Lane)

• Fix possible crash in array_position() or array_positions() when processing arrays of records (Junseok Yang)

• Fix one-byte buffer overrun in quote_literal_cstr() (Heikki Linnakangas)

  The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.

• Prevent multiple calls of pg_start_backup() and pg_stop_backup() from running concurrently (Michael Paquier)
This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

- Disable transform that attempted to remove no-op AT TIME ZONE conversions (Tom Lane)

  This resulted in wrong answers when the simplified expression was used in an index condition.

- Avoid discarding interval-to-interval casts that aren't really no-ops (Tom Lane)

  In some cases, a cast that should result in zeroing out low-order interval fields was mistakenly deemed to be a no-op and discarded. An example is that casting from INTERVAL MONTH to INTERVAL YEAR failed to clear the months field.

- Fix crash if the number of workers available to a parallel query decreases during a rescan (Andreas Seltenreich)

- Fix bugs in transmitting GUC parameter values to parallel workers (Michael Paquier, Tom Lane)

- Allow statements prepared with PREPARE to be given parallel plans (Amit Kapila, Tobias Bussmann)

- Fix incorrect generation of parallel plans for semi-joins (Tom Lane)

- Fix planner's cardinality estimates for parallel joins (Robert Haas)

  Ensure that these estimates reflect the number of rows predicted to be seen by each worker, rather than the total.

- Fix planner to avoid trying to parallelize plan nodes containing initplans or subplans (Tom Lane, Amit Kapila)

- Ensure that cached plans are invalidated by changes in foreign-table options (Amit Langote, Etsuro Fujita, Ashutosh Bapat)

- Fix the plan generated for sorted partial aggregation with a constant GROUP BY clause (Tom Lane)

- Fix “could not find plan for CTE” planner error when dealing with a UNION ALL containing CTE references (Tom Lane)

- Fix mishandling of initplans when forcibly adding a Material node to a subplan (Tom Lane)

  The typical consequence of this mistake was a “plan should not reference subplan's variable” error.

- Fix foreign-key-based join selectivity estimation for semi-joins and anti-joins, as well as inheritance cases (Tom Lane)

  The new code for taking the existence of a foreign key relationship into account did the wrong thing in these cases, making the estimates worse not better than the pre-9.6 code.

- Fix pg_dump to emit the data of a sequence that is marked as an extension configuration table (Michael Paquier)

- Fix mishandling of ALTER DEFAULT PRIVILEGES ... REVOKE in pg_dump (Stephen Frost)

  pg_dump missed issuing the required REVOKE commands in cases where ALTER DEFAULT PRIVILEGES had been used to reduce privileges to less than they would normally be.

- Fix pg_dump to dump user-defined casts and transforms that use built-in functions (Stephen Frost)

- Fix pg_restore with --create --if-exists to behave more sanely if an archive contains unrecognized DROP commands (Tom Lane)
This doesn't fix any live bug, but it may improve the behavior in future if `pg_restore` is used with an archive generated by a later `pg_dump` version.

- Fix `pg_basebackup`'s rate limiting in the presence of slow I/O (Antonin Houska)
  
  If disk I/O was transiently much slower than the specified rate limit, the calculation overflowed, effectively disabling the rate limit for the rest of the run.

- Fix `pg_basebackup`'s handling of symlinked `pg_stat_tmp` and `pg_replslot` subdirectories (Magnus Hagander, Michael Paquier)

- Fix possible `pg_basebackup` failure on standby server when including WAL files (Amit Kapila, Robert Haas)

- Improve `initdb` to insert the correct platform-specific default values for the `xxx_flush_after` parameters into `postgresql.conf` (Fabien Coelho, Tom Lane)
  
  This is a cleaner way of documenting the default values than was used previously.

- Fix possible mishandling of expanded arrays in domain check constraints and `CASE` execution (Tom Lane)
  
  It was possible for a PL/pgSQL function invoked in these contexts to modify or even delete an array value that needs to be preserved for additional operations.

- Fix nested uses of PL/pgSQL functions in contexts such as domain check constraints evaluated during assignment to a PL/pgSQL variable (Tom Lane)

- Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)
  
  This avoids failures if the objects are used after a Python garbage collection cycle has occurred.

- Fix PL/Tcl to support triggers on tables that have `.tupno` as a column name (Tom Lane)
  
  This matches the (previously undocumented) behavior of PL/Tcl's `spi_exec` and `spi_execp` commands, namely that a magic `.tupno` column is inserted only if there isn't a real column named that.

- Allow DOS-style line endings in `~/.pgpass` files, even on Unix (Vik Fearing)
  
  This change simplifies use of the same password file across Unix and Windows machines.

- Fix one-byte buffer overrun if `ecpg` is given a file name that ends with a dot (Takayuki Tsunakawa)

- Fix incorrect error reporting for duplicate data in `psql`'s `\crosstabview` (Tom Lane)
  
  `psql` sometimes quoted the wrong row and/or column values when complaining about multiple entries for the same crosstab cell.

- Fix `psql`'s tab completion for `ALTER DEFAULT PRIVILEGES` (Gilles Darold, Stephen Frost)

- Fix `psql`'s tab completion for `ALTER TABLE t ALTER c DROP` ... (Kyotaro Horiguchi)

- In `psql`, treat an empty or all-blank setting of the `PAGER` environment variable as meaning “no pager” (Tom Lane)
  
  Previously, such a setting caused output intended for the pager to vanish entirely.

- Improve `contrib/dblink`'s reporting of low-level libpq errors, such as out-of-memory (Joe Conway)
• Teach contrib/dblink to ignore irrelevant server options when it uses a contrib/postgres_fdw foreign server as the source of connection options (Corey Huinker)

Previously, if the foreign server object had options that were not also libpq connection options, an error occurred.

• Fix portability problems in contrib/pageinspect’s functions for GIN indexes (Peter Eisentraut, Tom Lane)

• Fix possible miss of socket read events while waiting on Windows (Amit Kapila)

This error was harmless for most uses, but it is known to cause hangs when trying to use the pldebugger extension.

• On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)

• Sync our copy of the timezone library with IANA release tzcode2016j (Tom Lane)

This fixes various issues, most notably that timezone data installation failed if the target directory didn’t support hard links.

• Update timezone data files to tzdata release 2016j for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarctica/Casey. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.

E.20. Release 9.6.1

Release date: 2016-10-27

This release contains a variety of fixes from 9.6.0. For information about new features in the 9.6 major release, see Section E.21.

E.20.1. Migration to Version 9.6.1

A dump/restore is not required for those running 9.6.X.

However, if your installation has been affected by the bugs described in the first two changelog entries below, then after updating you may need to take action to repair corrupted free space maps and/or visibility maps.

E.20.2. Changes

• Fix WAL-logging of truncation of relation free space maps and visibility maps (Pavan Deolasee, Heikki Linnakangas)

It was possible for these files to not be correctly restored during crash recovery, or to be written incorrectly on a standby server. Bogus entries in a free space map could lead to attempts to access pages that have been truncated away from the relation itself, typically producing errors like “could not read block XXX: read only 0 of 8192 bytes”. Checksum failures in the visibility map are also possible, if checksumming is enabled.

Procedures for determining whether there is a problem and repairing it if so are discussed at https://wiki.postgresql.org/wiki/Free_Space_Map_Problems.

• Fix possible data corruption when pg_upgrade rewrites a relation visibility map into 9.6 format (Tom Lane)
On big-endian machines, bytes of the new visibility map were written in the wrong order, leading to a completely incorrect map. On Windows, the old map was read using text mode, leading to incorrect results if the map happened to contain consecutive bytes that matched a carriage return/line feed sequence. The latter error would almost always lead to a pg_upgrade failure due to the map file appearing to be the wrong length.

If you are using a big-endian machine (many non-Intel architectures are big-endian) and have used pg_upgrade to upgrade from a pre-9.6 release, you should assume that all visibility maps are incorrect and need to be regenerated. It is sufficient to truncate each relation's visibility map with contrib/pg_visibility's pg_truncate_visibility_map() function. For more information see https://wiki.postgresql.org/wiki/Visibility_Map_Problems.

- Don’t throw serialization errors for self-conflicting insertions in \texttt{INSERT \ldots ON CONFLICT} (Thomas Munro, Peter Geoghegan)
- Fix use-after-free hazard in execution of aggregate functions using \texttt{DISTINCT} (Peter Geoghegan)

This could lead to a crash or incorrect query results.

- Fix incorrect handling of polymorphic aggregates used as window functions (Tom Lane)

The aggregate’s transition function was told that its first argument and result were of the aggregate’s output type, rather than the state type. This led to errors or crashes with polymorphic transition functions.

- Fix \texttt{COPY} with a column name list from a table that has row-level security enabled (Adam Brightwell)

- Fix \texttt{EXPLAIN} to emit valid XML when track_io_timing is on (Markus Winand)

Previously the XML output-format option produced syntactically invalid tags such as \texttt{<I/O-Read-Time>}. That is now rendered as \texttt{<I-O-Read-Time>}.

- Fix statistics update for \texttt{TRUNCATE} in a prepared transaction (Stas Kelvich)

- Fix bugs in merging inherited \texttt{CHECK} constraints while creating or altering a table (Tom Lane, Amit Langote)

Allow identical \texttt{CHECK} constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a \texttt{NOT VALID} constraint on the child. Likewise, prevent merging of a \texttt{NO INHERIT} child constraint with an inherited constraint.

- Show a sensible value in \texttt{pg_settings.unit} for \texttt{min_wal_size} and \texttt{max_wal_size} (Tom Lane)

- Fix replacement of array elements in \texttt{jsonb_set()} (Tom Lane)

If the target is an existing JSON array element, it got deleted instead of being replaced with a new value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)

- Preserve commit timestamps across server restart (Julien Rouhaud, Craig Ringer)

With track\_commit\_timestamp turned on, old commit timestamps became inaccessible after a clean server restart.

- Fix logical WAL decoding to work properly when a subtransaction’s WAL output is large enough to spill to disk (Andres Freund)

- Fix dangling-pointer problem in logical WAL decoding (Stas Kelvich)
Release Notes

• Round shared-memory allocation request to a multiple of the actual huge page size when attempting to use huge pages on Linux (Tom Lane)

   This avoids possible failures during `munmap()` on systems with atypical default huge page sizes. Except in crash-recovery cases, there were no ill effects other than a log message.

• Don't try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)

   This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)

   The reported problem involved leaking an error report during `PQreset()`, but there might be related cases.

• In `pg_upgrade`, check library loadability in name order (Tom Lane)

   This is a workaround to deal with cross-extension dependencies from language transform modules to their base language and data type modules.

• Fix `pg_upgrade` to work correctly for extensions containing index access methods (Tom Lane)

   To allow this, the server has been extended to support `ALTER EXTENSION ADD/DROP ACCESS METHOD`. That functionality should have been included in the original patch to support dynamic creation of access methods, but it was overlooked.

• Improve error reporting in `pg_upgrade`'s file copying/linking/rewriting steps (Tom Lane, Álvaro Herrera)

• Fix `pg_dump` to work against pre-7.4 servers (Amit Langote, Tom Lane)

• Disallow specifying both `--source-server` and `--source-target` options to `pg_rewind` (Michael Banck)

• Make `pg_rewind` turn off `synchronous_commit` in its session on the source server (Michael Banck, Michael Paquier)

   This allows `pg_rewind` to work even when the source server is using synchronous replication that is not working for some reason.

• In `pg_xlogdump`, retry opening new WAL segments when using `--follow` option (Magnus Hagander)

   This allows for a possible delay in the server's creation of the next segment.

• Fix `contrib/pg_visibility` to report the correct TID for a corrupt tuple that has been the subject of a rolled-back update (Tom Lane)

• Fix makefile dependencies so that parallel make of PL/Python by itself will succeed reliably (Pavel Raiskup)

• Update time zone data files to `tzdata` release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

   The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the `pg_timezone_names` view nor used for output.
In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.21. Release 9.6

Release date: 2016-09-29

E.21.1. Overview

Major enhancements in PostgreSQL 9.6 include:

• Parallel execution of sequential scans, joins and aggregates
• Avoid scanning pages unnecessarily during vacuum freeze operations
• Synchronous replication now allows multiple standby servers for increased reliability
• Full-text search can now search for phrases (multiple adjacent words)
• postgres_fdw now supports remote joins, sorts, UPDATEs, and DELETEs
• Substantial performance improvements, especially in the area of scalability on multi-CPU-socket servers

The above items are explained in more detail in the sections below.

E.21.2. Migration to Version 9.6

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.6 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

• Improve the pg_stat_activity view’s information about what a process is waiting for (Amit Kapila, Ildus Kurbangaliev)

Historically a process has only been shown as waiting if it was waiting for a heavyweight lock. Now waits for lightweight locks and buffer pins are also shown in pg_stat_activity. Also, the type of lock being waited for is now visible. These changes replace the waiting column with wait_event_type and wait_event.

• In to_char(), do not count a minus sign (when needed) as part of the field width for time-related fields (Bruce Momjian)

For example, to_char('−4 years'::interval, 'YY') now returns −04, rather than −4.

• Make extract() behave more reasonably with infinite inputs (Vitaly Burovoy)

Historically the extract() function just returned zero given an infinite timestamp, regardless of the given field name. Make it return infinity or −infinity as appropriate when the requested field is one that is monotonically increasing (e.g., year, epoch), or NULL when it is not (e.g., day, hour). Also, throw the expected error for bad field names.

• Remove PL/pgSQL’s “feature” that suppressed the innermost line of CONTEXT for messages emitted by RAISE commands (Pavel Stehule)

This ancient backwards-compatibility hack was agreed to have outlived its usefulness.

• Fix the default text search parser to allow leading digits in email and host tokens (Artur Zakirov)
In most cases this will result in few changes in the parsing of text. But if you have data where such addresses occur frequently, it may be worth rebuilding dependent tsvector columns and indexes so that addresses of this form will be found properly by text searches.

- Extend contrib/unaccent's standard unaccent.rules file to handle all diacritics known to Unicode, and to expand ligatures correctly (Thomas Munro, Léonard Benedetti)

  The previous version neglected to convert some less-common letters with diacritic marks. Also, ligatures are now expanded into separate letters. Installations that use this rules file may wish to rebuild tsvector columns and indexes that depend on the result.

- Remove the long-deprecated CREATEUSER/NOCREATEUSER options from CREATE ROLE and allied commands (Tom Lane)

  CREATEUSER actually meant SUPERUSER, for ancient backwards-compatibility reasons. This has been a constant source of confusion for people who (reasonably) expect it to mean CREATEROLE. It has been deprecated for ten years now, so fix the problem by removing it.

- Treat role names beginning with pg_ as reserved (Stephen Frost)

  User creation of such role names is now disallowed. This prevents conflicts with built-in roles created by initdb.

- Change a column name in the information_schema.routines view from result_cast_character_set_name to result_cast_char_set_name (Clément Prévost)

  The SQL:2011 standard specifies the longer name, but that appears to be a mistake, because adjacent column names use the shorter style, as do other information_schema views.

- psql's -c option no longer implies --no-psqlrc (Pavel Stehule, Catalin Iacob)

  Write --no-psqlrc (or its abbreviation -X) explicitly to obtain the old behavior. Scripts so modified will still work with old versions of psql.

- Improve pg_restore's -t option to match all types of relations, not only plain tables (Craig Ringer)

- Change the display format used for NextXID in pg_controldata and related places (Joe Conway, Bruce Momjian)

  Display epoch-and-transaction-ID values in the format number:number. The previous format number/number was confusingly similar to that used for LSNs.

- Update extension functions to be marked parallel-safe where appropriate (Andreas Karlsson)

  Many of the standard extensions have been updated to allow their functions to be executed within parallel query worker processes. These changes will not take effect in databases pg_upgrade'd from prior versions unless you apply ALTER EXTENSION UPDATE to each such extension (in each database of a cluster).

### E.21.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.6 and the previous major release.

#### E.21.3.1. Server

#### E.21.3.1.1. Parallel Queries

- Parallel queries (Robert Haas, Amit Kapila, David Rowley, many others)
With 9.6, PostgreSQL introduces initial support for parallel execution of large queries. Only strictly read-only queries where the driving table is accessed via a sequential scan can be parallelized. Hash joins and nested loops can be performed in parallel, as can aggregation (for supported aggregates). Much remains to be done, but this is already a useful set of features.

Parallel query execution is not (yet) enabled by default. To allow it, set the new configuration parameter max_parallel_workers_per_gather to a value larger than zero. Additional control over use of parallelism is available through other new configuration parameters force_parallel_mode, parallel_setup_cost, parallel_tuple_cost, and min_parallel_relation_size.

- Provide infrastructure for marking the parallel-safety status of functions (Robert Haas, Amit Kapila)

**E.21.3.1.2. Indexes**

- Allow GIN index builds to make effective use of maintenance_work_mem settings larger than 1 GB (Robert Abraham, Teodor Sigaev)
- Add pages deleted from a GIN index's pending list to the free space map immediately (Jeff Janes, Teodor Sigaev)

  This reduces bloat if the table is not vacuumed often.
- Add `gin_clean_pending_list()` function to allow manual invocation of pending-list cleanup for a GIN index (Jeff Janes)

  Formerly, such cleanup happened only as a byproduct of vacuuming or analyzing the parent table.
- Improve handling of dead index tuples in GiST indexes (Anastasia Lubennikova)

  Dead index tuples are now marked as such when an index scan notices that the corresponding heap tuple is dead. When inserting tuples, marked-dead tuples will be removed if needed to make space on the page.
- Add an SP-GiST operator class for type `box` (Alexander Lebedev)

**E.21.3.1.3. Sorting**

- Improve sorting performance by using quicksort, not replacement selection sort, when performing external sort steps (Peter Geoghegan)

  The new approach makes better use of the CPU cache for typical cache sizes and data volumes. Where necessary, the behavior can be adjusted via the new configuration parameter `replacement_sort_tuples`.
- Speed up text sorts where the same string occurs multiple times (Peter Geoghegan)
- Speed up sorting of `uuid`, `bytea`, and `char(n)` fields by using "abbreviated" keys (Peter Geoghegan)

  Support for abbreviated keys has also been added to the non-default operator classes `text_pattern_ops`, `varchar_pattern_ops`, and `bpchar_pattern_ops`. Processing of ordered-set aggregates can also now exploit abbreviated keys.
- Speed up `CREATE INDEX CONCURRENTLY` by treating TIDs as 64-bit integers during sorting (Peter Geoghegan)

**E.21.3.1.4. Locking**

- Reduce contention for the `ProcArrayLock` (Amit Kapila, Robert Haas)
• Improve performance by moving buffer content locks into the buffer descriptors (Andres Freund, Simon Riggs)

• Replace shared-buffer header spinlocks with atomic operations to improve scalability (Alexander Korotkov, Andres Freund)

• Use atomic operations, rather than a spinlock, to protect an LWLock's wait queue (Andres Freund)

• Partition the shared hash table freelist to reduce contention on multi-CPU-socket servers (Alexander Alekseev)

• Reduce interlocking on standby servers during the replay of btree index vacuuming operations (Simon Riggs)

This change avoids substantial replication delays that sometimes occurred while replaying such operations.

E.21.3.1.5. Optimizer Statistics

• Improve ANALYZE's estimates for columns with many nulls (Tomas Vondra, Alex Shulgin)

Previously ANALYZE tended to underestimate the number of non-NULL distinct values in a column with many NULLs, and was also inaccurate in computing the most-common values.

• Improve planner's estimate of the number of distinct values in a query result (Tomas Vondra)

• Use foreign key relationships to infer selectivity for join predicates (Tomas Vondra, David Rowley)

If a table t has a foreign key restriction, say (a, b) REFERENCES r (x, y), then a WHERE condition such as t.a = r.x AND t.b = r.y cannot select more than one r row per t row. The planner formerly considered these AND conditions to be independent and would often drastically misestimate selectivity as a result. Now it compares the WHERE conditions to applicable foreign key constraints and produces better estimates.

E.21.3.1.6. VACUUM

• Avoid re-vacuuming pages containing only frozen tuples (Masahiko Sawada, Robert Haas, Andres Freund)

Formerly, anti-wraparound vacuum had to visit every page of a table, even pages where there was nothing to do. Now, pages containing only already-frozen tuples are identified in the table's visibility map, and can be skipped by vacuum even when doing transaction wraparound prevention. This should greatly reduce the cost of maintaining large tables containing mostly-unchanging data.

If necessary, vacuum can be forced to process all-frozen pages using the new DISABLE_PAGE_SKIPPING option. Normally this should never be needed, but it might help in recovering from visibility-map corruption.

• Avoid useless heap-truncation attempts during VACUUM (Jeff Janes, Tom Lane)

This change avoids taking an exclusive table lock in some cases where no truncation is possible. The main benefit comes from avoiding unnecessary query cancellations on standby servers.

E.21.3.1.7. General Performance

• Allow old MVCC snapshots to be invalidated after a configurable timeout (Kevin Grittner)

Normally, deleted tuples cannot be physically removed by vacuuming until the last transaction that could “see” them is gone. A transaction that stays open for a long time can thus cause considerable table bloat because space cannot be recycled. This feature allows setting a time-based limit, via the new configuration parameter old_snapshot_threshold, on how long an MVCC snapshot is guaran-
ted to be valid. After that, dead tuples are candidates for removal. A transaction using an outdated snapshot will get an error if it attempts to read a page that potentially could have contained such data.

- Ignore `GROUP BY` columns that are functionally dependent on other columns (David Rowley)

If a `GROUP BY` clause includes all columns of a non-deferred primary key, as well as other columns of the same table, those other columns are redundant and can be dropped from the grouping. This saves computation in many common cases.

- Allow use of an index-only scan on a partial index when the index's `WHERE` clause references columns that are not indexed (Tomas Vondra, Kyotaro Horiguchi)

For example, an index defined by `CREATE INDEX tidx_partial ON t(b) WHERE a > 0` can now be used for an index-only scan by a query that specifies `WHERE a > 0` and does not otherwise use `a`. Previously this was disallowed because `a` is not listed as an index column.

- Perform checkpoint writes in sorted order (Fabien Coelho, Andres Freund)

Previously, checkpoints wrote out dirty pages in whatever order they happen to appear in shared buffers, which usually is nearly random. That performs poorly, especially on rotating media. This change causes checkpoint-driven writes to be done in order by file and block number, and to be balanced across tablespaces.

- Where feasible, trigger kernel writeback after a configurable number of writes, to prevent accumulation of dirty data in kernel disk buffers (Fabien Coelho, Andres Freund)

PostgreSQL writes data to the kernel's disk cache, from where it will be flushed to physical storage in due time. Many operating systems are not smart about managing this and allow large amounts of dirty data to accumulate before deciding to flush it all at once, causing long delays for new I/O requests until the flushing finishes. This change attempts to alleviate this problem by explicitly requesting data flushes after a configurable interval.

On Linux, `sync_file_range()` is used for this purpose, and the feature is on by default on Linux because that function has few downsides. This flushing capability is also available on other platforms if they have `msync()` or `posix_fadvise()`, but those interfaces have some undesirable side-effects so the feature is disabled by default on non-Linux platforms.

The new configuration parameters `backend_flush_after`, `bgwriter_flush_after`, `checkpoint_flush_after`, and `wal_writer_flush_after` control this behavior.

- Improve aggregate-function performance by sharing calculations across multiple aggregates if they have the same arguments and transition functions (David Rowley)

For example, `SELECT AVG(x), VARIANCE(x) FROM tab` can use a single per-row computation for both aggregates.

- Speed up visibility tests for recently-created tuples by checking the current transaction's snapshot, not `pg_clog`, to decide if the source transaction should be considered committed (Jeff Janes, Tom Lane)

- Allow tuple hint bits to be set sooner than before (Andres Freund)

- Improve performance of short-lived prepared transactions (Stas Kelvich, Simon Riggs, Pavan Deolasee)

Two-phase commit information is now written only to WAL during `PREPARE TRANSACTION`, and will be read back from WAL during `COMMIT PREPARED` if that happens soon thereafter. A separate state file is created only if the pending transaction does not get committed or aborted by the time of the next checkpoint.

- Improve performance of memory context destruction (Jan Wieck)
Release Notes

- Improve performance of resource owners with many tracked objects (Aleksander Alekseev)
- Improve speed of the output functions for timestamp, time, and date data types (David Rowley, Andres Freund)
- Avoid some unnecessary cancellations of hot-standby queries during replay of actions that take AccessExclusive locks (Jeff Janes)
- Extend relations multiple blocks at a time when there is contention for the relation's extension lock (Dilip Kumar)

This improves scalability by decreasing contention.
- Increase the number of clog buffers for better scalability (Amit Kapila, Andres Freund)
- Speed up expression evaluation in PL/pgSQL by keeping ParamListInfo entries for simple variables valid at all times (Tom Lane)
- Avoid reducing the SO_SNDBUF setting below its default on recent Windows versions (Chen Hua-jun)
- Disable update_process_title by default on Windows (Takayuki Tsunakawa)

The overhead of updating the process title is much larger on Windows than most other platforms, and it is also less useful to do it since most Windows users do not have tools that can display process titles.

E.21.3.1.8. Monitoring

- Add pg_stat_progress_vacuum system view to provide progress reporting for VACUUM operations (Amit Langote, Robert Haas, Vinayak Pokale, Rahila Syed)
- Add pg_control_system(), pg_control_checkpoint(), pg_control_recovery(), and pg_control_init() functions to expose fields of pg_control to SQL (Joe Conway, Michael Paquier)
- Add pg_config system view (Joe Conway)

This view exposes the same information available from the pg_config command-line utility, namely assorted compile-time configuration information for PostgreSQL.
- Add a confirmed_flush_lsn column to the pg_replication_slots system view (Marko Tiikkaja)
- Add pg_stat_wal_receiver system view to provide information about the state of a hot-standby server's WAL receiver process (Michael Paquier)
- Add pg_blocking_pids() function to reliably identify which sessions block which others (Tom Lane)

This function returns an array of the process IDs of any sessions that are blocking the session with the given process ID. Historically users have obtained such information using a self-join on the pg_locks view. However, it is unreasonably tedious to do it that way with any modicum of correctness, and the addition of parallel queries has made the old approach entirely impractical, since locks might be held or awaited by child worker processes rather than the session's main process.
- Add function pg_current_xlog_flush_location() to expose the current transaction log flush location (Tomas Vondra)
- Add function pg_notification_queue_usage() to report how full the NOTIFY queue is (Brendan Jurd)
• Limit the verbosity of memory context statistics dumps (Tom Lane)

The memory usage dump that is output to the postmaster log during an out-of-memory failure now summarizes statistics when there are a large number of memory contexts, rather than possibly generating a very large report. There is also a “grand total” summary line now.

E.21.3.1.9. Authentication

• Add a BSD authentication method to allow use of the BSD Authentication service for PostgreSQL client authentication (Marisa Emerson)

BSD Authentication is currently only available on OpenBSD.

• When using PAM authentication, provide the client IP address or host name to PAM modules via the PAM_RHOST item (Grzegorz Sampolski)

• Provide detail in the postmaster log for more types of password authentication failure (Tom Lane)

All ordinarily-reachable password authentication failure cases should now provide specific DETAIL fields in the log.

• Support RADIUS passwords up to 128 characters long (Marko Tiikkaja)

• Add new SSPI authentication parameters compat_realm and upn_username to control whether NetBIOS or Kerberos realm names and user names are used during SSPI authentication (Christian Ullrich)

E.21.3.1.10. Server Configuration

• Allow sessions to be terminated automatically if they are in idle-in-transaction state for too long (Vik Fearing)

This behavior is controlled by the new configuration parameter idle_in_transaction_session_timeout. It can be useful to prevent forgotten transactions from holding locks or preventing vacuum cleanup for too long.

• Raise the maximum allowed value of checkpoint_timeout to 24 hours (Simon Riggs)

• Allow effective_io_concurrency to be set per-tablespace to support cases where different tablespaces have different I/O characteristics (Julien Rouhaud)

• Add log_line_prefix option %n to print the current time in Unix epoch form, with milliseconds (Tomas Vondra, Jeff Davis)

• Add syslog_sequence_numbers and syslog_split_messages configuration parameters to provide more control over the message format when logging to syslog (Peter Eisentraut)

• Merge the archive and hot_standby values of the wal_level configuration parameter into a single new value replica (Peter Eisentraut)

Making a distinction between these settings is no longer useful, and merging them is a step towards a planned future simplification of replication setup. The old names are still accepted but are converted to replica internally.

• Add configure option --with-systemd to enable calling sd_notify() at server start and stop (Peter Eisentraut)

This allows the use of systemd service units of type notify, which greatly simplifies the management of PostgreSQL under systemd.

• Allow the server's SSL key file to have group read access if it is owned by root (Christoph Berg)
Formerly, we insisted the key file be owned by the user running the PostgreSQL server, but that is inconvenient on some systems (such as Debian) that are configured to manage certificates centrally. Therefore, allow the case where the key file is owned by root and has group read access. It is up to the operating system administrator to ensure that the group does not include any untrusted users.

**E.21.3.11. Reliability**

- Force backends to exit if the postmaster dies (Rajeev Rastogi, Robert Haas)

  Under normal circumstances the postmaster should always outlive its child processes. If for some reason the postmaster dies, force backend sessions to exit with an error. Formerly, existing backends would continue to run until their clients disconnect, but that is unsafe and inefficient. It also prevents a new postmaster from being started until the last old backend has exited. Backends will detect postmaster death when waiting for client I/O, so the exit will not be instantaneous, but it should happen no later than the end of the current query.

- Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

  When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

- Ensure that invalidation messages are recorded in WAL even when issued by a transaction that has no XID assigned (Andres Freund)

  This fixes some corner cases in which transactions on standby servers failed to notice changes, such as new indexes.

- Prevent multiple processes from trying to clean a GIN index's pending list concurrently (Teodor Sigaev, Jeff Janes)

  This had been intentionally allowed, but it causes race conditions that can result in vacuum missing index entries it needs to delete.

**E.21.3.2. Replication and Recovery**

- Allow synchronous replication to support multiple simultaneous synchronous standby servers, not just one (Masahiko Sawada, Beena Emerson, Michael Paquier, Fujii Masao, Kyotaro Horiguchi)

  The number of standby servers that must acknowledge a commit before it is considered complete is now configurable as part of the synchronous_standby_names parameter.

- Add new setting `remote_apply` for configuration parameter synchronous_commit (Thomas Munro)

  In this mode, the master waits for the transaction to be applied on the standby server, not just written to disk. That means that you can count on a transaction started on the standby to see all commits previously acknowledged by the master.

- Add a feature to the replication protocol, and a corresponding option to `pg_create_physical_replication_slot()`, to allow reserving WAL immediately when creating a replication slot (Gurjeet Singh, Michael Paquier)

  This allows the creation of a replication slot to guarantee that all the WAL needed for a base backup will be available.

- Add a `--slot` option to `pg_basebackup` (Peter Eisentraut)
This lets pg_basebackup use a replication slot defined for WAL streaming. After the base backup completes, selecting the same slot for regular streaming replication allows seamless startup of the new standby server.

- Extend `pg_start_backup()` and `pg_stop_backup()` to support non-exclusive backups (Magnus Hagander)

### E.21.3.3. Queries

- Allow functions that return sets of tuples to return simple NULLs (Andrew Gierth, Tom Lane)

  In the context of `SELECT FROM function(...)`, a function that returned a set of composite values was previously not allowed to return a plain NULL value as part of the set. Now that is allowed and interpreted as a row of NULLs. This avoids corner-case errors with, for example, unnesting an array of composite values.

- Fully support array subscripts and field selections in the target column list of an `INSERT` with multiple `VALUES` rows (Tom Lane)

  Previously, such cases failed if the same target column was mentioned more than once, e.g.,
  ```sql
  INSERT INTO tab (x[1], x[2]) VALUES (...).
  ```

- When appropriate, postpone evaluation of `SELECT` output expressions until after an `ORDER BY` sort (Konstantin Knizhnik)

  This change ensures that volatile or expensive functions in the output list are executed in the order suggested by `ORDER BY`, and that they are not evaluated more times than required when there is a `LIMIT` clause. Previously, these properties held if the ordering was performed by an index scan or pre-merge-join sort, but not if it was performed by a top-level sort.

- Widen counters recording the number of tuples processed to 64 bits (Andreas Scherbaum)

  This change allows command tags, e.g. `SELECT`, to correctly report tuple counts larger than 4 billion. This also applies to PL/pgSQL's `GET DIAGNOSTICS ... ROW_COUNT` command.

- Avoid doing encoding conversions by converting through the MULE_INTERNAL encoding (Tom Lane)

  Previously, many conversions for Cyrillic and Central European single-byte encodings were done by converting to a related MULE_INTERNAL coding scheme and then to the destination encoding. Aside from being inefficient, this meant that when the conversion encountered an untranslatable character, the error message would confusingly complain about failure to convert to or from MULE_INTERNAL, rather than the user-visible encoding.

- Consider performing joins of foreign tables remotely only when the tables will be accessed under the same role ID (Shigeru Hanada, Ashutosh Bapat, Etsuro Fujita)

  Previously, the foreign join pushdown infrastructure left the question of security entirely up to individual foreign data wrappers, but that made it too easy for an FDW to inadvertently create subtle security holes. So, make it the core code's job to determine which role ID will access each table, and do not attempt join pushdown unless the role is the same for all relevant relations.

### E.21.3.4. Utility Commands

- Allow `COPY` to copy the output of an `INSERT/UPDATE/DELETE ... RETURNING` query (Marko Tiikkaja)

  Previously, an intermediate CTE had to be written to get this result.

- Introduce `ALTER object DEPENDS ON EXTENSION` (Abhijit Menon-Sen)
This command allows a database object to be marked as depending on an extension, so that it will be dropped automatically if the extension is dropped (without needing \texttt{CASCADE}). However, the object is not part of the extension, and thus will be dumped separately by \texttt{pg_dump}.

- Make \texttt{ALTER \_object \_SET \_SCHEMA} do nothing when the object is already in the requested schema, rather than throwing an error as it historically has for most object types (Marti Raudsepp)

- Add options to \texttt{ALTER \_OPERATOR} to allow changing the selectivity functions associated with an existing operator (Yury Zhuravlev)

- Add an \texttt{IF \_NOT \_EXISTS} option to \texttt{ALTER \_TABLE \_ADD \_COLUMN} (Fabrízio de Royes Mello)

- Reduce the lock strength needed by \texttt{ALTER \_TABLE} when setting fillfactor and autovacuum-related relation options (Fabrízio de Royes Mello, Simon Riggs)

- Introduce \texttt{CREATE \_ACCESS \_METHOD} to allow extensions to create index access methods (Alexander Korotkov, Petr Jelinek)

- Add a \texttt{CASCADE} option to \texttt{CREATE \_EXTENSION} to automatically create any extensions the requested one depends on (Petr Jelinek)

- Make \texttt{CREATE \_TABLE \_\_LIKE} include an \texttt{OID} column if any source table has one (Bruce Momjian)

- If a \texttt{CHECK} constraint is declared \texttt{NOT \_VALID} in a table creation command, automatically mark it as valid (Amit Langote, Amul Sul)

  This is safe because the table has no existing rows. This matches the longstanding behavior of \texttt{FOREIGN \_KEY} constraints.

- Fix \texttt{DROP \_OPERATOR} to clear \texttt{pg\_operator\_oprcom} and \texttt{pg\_operator\_oprnegate} links to the dropped operator (Roma Sokolov)

  Formerly such links were left as-is, which could pose a problem in the somewhat unlikely event that the dropped operator's \texttt{OID} was reused for another operator.

- Do not show the same subplan twice in \texttt{EXPLAIN} output (Tom Lane)

  In certain cases, typically involving SubPlan nodes in index conditions, \texttt{EXPLAIN} would print data for the same subplan twice.

- Disallow creation of indexes on system columns, except for \texttt{OID} columns (David Rowley)

  Such indexes were never considered supported, and would very possibly misbehave since the system might change the system-column fields of a tuple without updating indexes. However, previously there were no error checks to prevent them from being created.

\textbf{E.21.3.5. Permissions Management}

- Use the privilege system to manage access to sensitive functions (Stephen Frost)

  Formerly, many security-sensitive functions contained hard-wired checks that would throw an error if they were called by a non-superuser. This forced the use of superuser roles for some relatively pedestrian tasks. The hard-wired error checks are now gone in favor of making \texttt{initdb} revoke the default public \texttt{EXECUTE} privilege on these functions. This allows installations to choose to grant usage of such functions to trusted roles that do not need all superuser privileges.

- Create some built-in roles that can be used to grant access to what were previously superuser-only functions (Stephen Frost)
Currently the only such role is `pg_signal_backend`, but more are expected to be added in future.

### E.21.3.6. Data Types

- Improve full-text search to support searching for phrases, that is, lexemes appearing adjacent to each other in a specific order, or with a specified distance between them (Teodor Sigaev, Oleg Bartunov, Dmitry Ivanov)

  A phrase-search query can be specified in `tsquery` input using the new operators `<->` and `<N>`. The former means that the lexemes before and after it must appear adjacent to each other in that order. The latter means they must be exactly `N` lexemes apart.

- Allow omitting one or both boundaries in an array slice specifier, e.g. `array_col[3:]` (Yury Zhuravlev)

  Omitted boundaries are taken as the upper or lower limit of the corresponding array subscript. This allows simpler specification for many common use-cases.

- Be more careful about out-of-range dates and timestamps (Vitaly Burovoy)

  This change prevents unexpected out-of-range errors for `timestamp with time zone` values very close to the implementation limits. Previously, the “same” value might be accepted or not depending on the `timezone` setting, meaning that a dump and reload could fail on a value that had been accepted when presented. Now the limits are enforced according to the equivalent UTC time, not local time, so as to be independent of `timezone`.

  Also, PostgreSQL is now more careful to detect overflow in operations that compute new date or timestamp values, such as `date + integer`.

- For geometric data types, make sure `infinity` and `NaN` component values are treated consistently during input and output (Tom Lane)

  Such values will now always print the same as they would in a simple `float8` column, and be accepted the same way on input. Previously the behavior was platform-dependent.

- Upgrade the `ispell` dictionary type to handle modern Hunspell files and support more languages (Artur Zakirov)

- Implement look-behind constraints in regular expressions (Tom Lane)

  A look-behind constraint is like a lookahead constraint in that it consumes no text; but it checks for existence (or nonexistence) of a match ending at the current point in the string, rather than one starting at the current point. Similar features exist in many other regular-expression engines.

- In regular expressions, if an apparent three-digit octal escape `\nnn` would exceed 377 (255 decimal), assume it is a two-digit octal escape instead (Tom Lane)

  This makes the behavior match current Tcl releases.

- Add transaction ID operators `xid <> xid` and `xid <> int4`, for consistency with the corresponding equality operators (Michael Paquier)

### E.21.3.7. Functions

- Add `jsonb_insert()` function to insert a new element into a `jsonb` array, or a not-previous-ly-existing key into a `jsonb` object (Dmitry Dolgov)

- Improve the accuracy of the `ln()`, `log()`, `exp()`, and `pow()` functions for type `numeric` (Dean Rasheed)
• Add a `scale(numeric)` function to extract the display scale of a numeric value (Marko Tiikkaja)

• Add trigonometric functions that work in degrees (Dean Rasheed)

For example, `sind()` measures its argument in degrees, whereas `sin()` measures in radians. These functions go to some lengths to deliver exact results for values where an exact result can be expected, for instance `sind(30) = 0.5`.

• Ensure that trigonometric functions handle infinity and NaN inputs per the POSIX standard (Dean Rasheed)

The POSIX standard says that these functions should return NaN for NaN input, and should throw an error for out-of-range inputs including infinity. Previously our behavior varied across platforms.

• Make `to_timestamp(float8)` convert float infinity to timestamp infinity (Vitaly Burovoy)

Formerly it just failed on an infinite input.

• Add new functions for `tsvector` data (Stas Kelvich)

The new functions are `ts_delete()`, `ts_filter()`, `unnest()`, `tsvector_to_array()`, `array_to_tsvector()`, and a variant of `setweight()` that sets the weight only for specified lexeme(s).

• Allow `ts_stat()` and `tsvector_update_trigger()` to operate on values that are of types binary-compatible with the expected argument type, not just exactly that type; for example allow `citext` where `text` is expected (Teodor Sigaev)

• Add variadic functions `num_nulls()` and `num_nonnulls()` that count the number of their arguments that are null or non-null (Marko Tiikkaja)

An example usage is `CHECK(num_nonnulls(a,b,c) = 1)` which asserts that exactly one of `a,b,c` is not `NULL`. These functions can also be used to count the number of null or nonnull elements in an array.

• Add function `parse_ident()` to split a qualified, possibly quoted SQL identifier into its parts (Pavel Stehule)

• In `to_number()`, interpret a `V` format code as dividing by 10 to the power of the number of digits following `V` (Bruce Momjian)

This makes it operate in an inverse fashion to `to_char()`.

• Make the `to_reg*()` functions accept type `text` not `cstring` (Petr Korobeinikov)

This avoids the need to write an explicit cast in most cases where the argument is not a simple literal constant.

• Add `pg_size_bytes()` function to convert human-readable size strings to numbers (Pavel Stehule, Vitaly Burovoy, Dean Rasheed)

This function converts strings like those produced by `pg_size_pretty()` into bytes. An example usage is `SELECT oid::regclass FROM pg_class WHERE pg_total_relation_size(oid) > pg_size_bytes('10 GB')`.

• In `pg_size_pretty()`, format negative numbers similarly to positive ones (Adrian Vondendriesch)

Previously, negative numbers were never abbreviated, just printed in bytes.
• Add an optional missing_ok argument to the current_setting() function (David Christieensen)

This allows avoiding an error for an unrecognized parameter name, instead returning a NULL.

• Change various catalog-inspection functions to return NULL for invalid input (Michael Paquier)

pg_get_viewdef() now returns NULL if given an invalid view OID, and several similar functions likewise return NULL for bad input. Previously, such cases usually led to “cache lookup failed” errors, which are not meant to occur in user-facing cases.

• Fix pg_replication_origin_xact_reset() to not have any arguments (Fujii Masao)

The documentation said that it has no arguments, and the C code did not expect any arguments, but the entry in pg_proc mistakenly specified two arguments.

E.21.3.8. Server-Side Languages

• In PL/pgSQL, detect mismatched CONTINUE and EXIT statements while compiling a function, rather than at execution time (Jim Nasby)

• Extend PL/Python’s error-reporting and message-reporting functions to allow specifying additional message fields besides the primary error message (Pavel Stehule)

• Allow PL/Python functions to call themselves recursively via SPI, and fix the behavior when multiple set-returning PL/Python functions are called within one query (Alexey Grishchenko, Tom Lane)

• Fix session-lifespan memory leaks in PL/Python (Heikki Linnakangas, Haribabu Kommi, Tom Lane)

• Modernize PL/Tcl to use Tcl’s “object” APIs instead of simple strings (Jim Nasby, Karl Lehenbauer)

This can improve performance substantially in some cases. Note that PL/Tcl now requires Tcl 8.4 or later.

• In PL/Tcl, make database-reported errors return additional information in Tcl's errorCode global variable (Jim Nasby, Tom Lane)

This feature follows the Tcl convention for returning auxiliary data about an error.

• Fix PL/Tcl to perform encoding conversion between the database encoding and UTF-8, which is what Tcl expects (Tom Lane)

Previously, strings were passed through without conversion, leading to misbehavior with non-ASCII characters when the database encoding was not UTF-8.

E.21.3.9. Client Interfaces

• Add a nonlocalized version of the severity field in error and notice messages (Tom Lane)

This change allows client code to determine severity of an error or notice without having to worry about localized variants of the severity strings.

• Introduce a feature in libpq whereby the CONTEXT field of messages can be suppressed, either always or only for non-error messages (Pavel Stehule)

The default behavior of PQerrorMessage() is now to print CONTEXT only for errors. The new function PQsetErrorContextVisibility() can be used to adjust this.

• Add support in libpq for regenerating an error message with a different verbosity level (Alex Shulgin)
This is done with the new function PQresultVerboseErrorMessage(). This supports psql's new \errverbose feature, and may be useful for other clients as well.

- Improve libpq's PQhost() function to return useful data for default Unix-socket connections (Tom Lane)

  Previously it would return NULL if no explicit host specification had been given; now it returns the default socket directory path.

- Fix ecpg's lexer to handle line breaks within comments starting on preprocessor directive lines (Michael Meskes)

### E.21.3.10. Client Applications

- Add a --strict-names option to pg_dump and pg_restore (Pavel Stehule)

  This option causes the program to complain if there is no match for a -t or -n option, rather than silently doing nothing.

- In pg_dump, dump locally-made changes of privilege assignments for system objects (Stephen Frost)

  While it has always been possible for a superuser to change the privilege assignments for built-in or extension-created objects, such changes were formerly lost in a dump and reload. Now, pg_dump recognizes and dumps such changes. (This works only when dumping from a 9.6 or later server, however.)

- Allow pg_dump to dump non-extension-owned objects that are within an extension-owned schema (Martín Marqués)

  Previously such objects were ignored because they were mistakenly assumed to belong to the extension owning their schema.

- In pg_dump output, include the table name in object tags for object types that are only uniquely named per-table (for example, triggers) (Peter Eisentraut)

### E.21.3.10.1. psql

- Support multiple --c and --f command-line options (Pavel Stehule, Catalin Iacob)

  The specified operations are carried out in the order in which the options are given, and then psql terminates.

- Add a \crosstabview command that prints the results of a query in a cross-tabulated display (Daniel Vérité)

  In the crosstab display, data values from one query result column are placed in a grid whose column and row headers come from other query result columns.

- Add an \errverbose command that shows the last server error at full verbosity (Alex Shulgin)

  This is useful after getting an unexpected error — you no longer need to adjust the VERBOSITY variable and recreate the failure in order to see error fields that are not shown by default.

- Add \ev and \sv commands for editing and showing view definitions (Petr Korobeinikov)

  These are parallel to the existing \ef and \sf commands for functions.

- Add a \gexec command that executes a query and re-submits the result(s) as new queries (Corey Huinker)
Release Notes

• Allow \pset C string to set the table title, for consistency with \C string (Bruce Momjian)

• In \pset expanded auto mode, do not use expanded format for query results with only one column (Andreas Karlsson, Robert Haas)

• Improve the headers output by the \watch command (Michael Paquier, Tom Lane)

  Include the \pset title string if one has been set, and shorten the prefabricated part of the header to be timestamp (every Ns). Also, the timestamp format now obeys psql's locale environment.

• Improve tab-completion logic to consider the entire input query, not only the current line (Tom Lane)

  Previously, breaking a command into multiple lines defeated any tab completion rules that needed to see words on earlier lines.

• Numerous minor improvements in tab-completion behavior (Peter Eisentraut, Vik Fearing, Kevin Grittner, Kyotaro Horiguchi, Jeff Janes, Andreas Karlsson, Fujii Masao, Thomas Munro, Masahiko Sawada, Pavel Stehule)

• Add a PROMPT option %p to insert the process ID of the connected backend (Julien Rouhaud)

• Introduce a feature whereby the CONTEXT field of messages can be suppressed, either always or only for non-error messages (Pavel Stehule)

  Printing CONTEXT only for errors is now the default behavior. This can be changed by setting the special variable SHOW_CONTEXT.

• Make \df+ show function access privileges and parallel-safety attributes (Michael Paquier)

E.21.3.10.2. pgbench

• SQL commands in pgbench scripts are now ended by semicolons, not newlines (Kyotaro Horiguchi, Tom Lane)

  This change allows SQL commands in scripts to span multiple lines. Existing custom scripts will need to be modified to add a semicolon at the end of each line that does not have one already. (Doing so does not break the script for use with older versions of pgbench.)

• Support floating-point arithmetic, as well as some built-in functions, in expressions in backslash commands (Fabien Coelho)

• Replace \setrandom with built-in functions (Fabien Coelho)

  The new built-in functions include random(), random_exponential(), and random_gaussian(), which perform the same work as \setrandom, but are easier to use since they can be embedded in larger expressions. Since these additions have made \setrandom obsolete, remove it.

• Allow invocation of multiple copies of the built-in scripts, not only custom scripts (Fabien Coelho)

  This is done with the new -b switch, which works similarly to -f for custom scripts.

• Allow changing the selection probabilities (weights) for scripts (Fabien Coelho)

  When multiple scripts are specified, each pgbench transaction randomly chooses one to execute. Formerly this was always done with uniform probability, but now different selection probabilities can be specified for different scripts.

• Collect statistics for each script in a multi-script run (Fabien Coelho)
This feature adds an intermediate level of detail to existing global and per-command statistics print-outs.

• Add a \texttt{--progress-timestamp} option to report progress with Unix epoch timestamps, instead of time since the run started (Fabien Coelho)

• Allow the number of client connections (\texttt{-c}) to not be an exact multiple of the number of threads (\texttt{-j}) (Fabien Coelho)

• When the \texttt{-T} option is used, stop promptly at the end of the specified time (Fabien Coelho)

Previously, specifying a low transaction rate could cause pgbench to wait significantly longer than specified.

\subsection*{E.21.3.11. Server Applications}

• Improve error reporting during initdb's post-bootstrap phase (Tom Lane)

Previously, an error here led to reporting the entire input file as the “failing query”; now just the current query is reported. To get the desired behavior, queries in initdb's input files must be separated by blank lines.

• Speed up initdb by using just one standalone-backend session for all the post-bootstrap steps (Tom Lane)

• Improve pg_rewind so that it can work when the target timeline changes (Alexander Korotkov)

This allows, for example, rewinding a promoted standby back to some state of the old master's timeline.

\subsection*{E.21.3.12. Source Code}

• Remove obsolete \texttt{heap_formtuple/heap_modifytuple/heap_deformtuple} functions (Peter Geoghegan)

• Add macros to make \texttt{AllocSetContextCreate()} calls simpler and safer (Tom Lane)

Writing out the individual sizing parameters for a memory context is now deprecated in favor of using one of the new macros \texttt{ALLOCSET_DEFAULT_SIZES}, \texttt{ALLOCSET_SMALL_SIZES}, or \texttt{ALLOCSET_START_SMALL_SIZES}. Existing code continues to work, however.

• Unconditionally use static inline functions in header files (Andres Freund)

This may result in warnings and/or wasted code space with very old compilers, but the notational improvement seems worth it.

• Improve TAP testing infrastructure (Michael Paquier, Craig Ringer, Álvaro Herrera, Stephen Frost)

Notably, it is now possible to test recovery scenarios using this infrastructure.

• Make \texttt{trace_lwlocks} identify individual locks by name (Robert Haas)

• Improve \texttt{psql}'s tab-completion code infrastructure (Thomas Munro, Michael Paquier)

Tab-completion rules are now considerably easier to write, and more compact.

• Nail the \texttt{pg_shseclabel} system catalog into cache, so that it is available for access during connection authentication (Adam Brightwell)

The core code does not use this catalog for authentication, but extensions might wish to consult it.
• Restructure index access method API to hide most of it at the C level (Alexander Korotkov, Andrew Gierth)

This change modernizes the index AM API to look more like the designs we have adopted for foreign data wrappers and tablesample handlers. This simplifies the C code and makes it much more practical to define index access methods in installable extensions. A consequence is that most of the columns of the \texttt{pg\_am} system catalog have disappeared. New inspection functions have been added to allow SQL queries to determine index AM properties that used to be discoverable from \texttt{pg\_am}.

• Add \texttt{pg\_init\_privs} system catalog to hold original privileges of initdb-created and extension-created objects (Stephen Frost)

This infrastructure allows \texttt{pg\_dump} to dump changes that an installation may have made in privileges attached to system objects. Formerly, such changes would be lost in a dump and reload, but now they are preserved.

• Change the way that extensions allocate custom LWLocks (Amit Kapila, Robert Haas)

The \texttt{RequestAddinLWLocks()} function is removed, and replaced by \texttt{RequestNamedLWLockTranche()}. This allows better identification of custom LWLocks, and is less error-prone.

• Improve the isolation tester to allow multiple sessions to wait concurrently, allowing testing of deadlock scenarios (Robert Haas)

• Introduce extensible node types (KaiGai Kohei)

This change allows FDWs or custom scan providers to store data in a plan tree in a more convenient format than was previously possible.

• Make the planner deal with post-scan/join query steps by generating and comparing Paths, replacing a lot of ad-hoc logic (Tom Lane)

This change provides only marginal user-visible improvements today, but it enables future work on a lot of upper-planner improvements that were impractical to tackle using the old code structure.

• Support partial aggregation (David Rowley, Simon Riggs)

This change allows the computation of an aggregate function to be split into separate parts, for example so that parallel worker processes can cooperate on computing an aggregate. In future it might allow aggregation across local and remote data to occur partially on the remote end.

• Add a generic command progress reporting facility (Vinayak Pokale, Rahila Syed, Amit Langote, Robert Haas)

• Separate out psql's flex lexer to make it usable by other client programs (Tom Lane, Kyotaro Horiguchi)

This eliminates code duplication for programs that need to be able to parse SQL commands well enough to identify command boundaries. Doing that in full generality is more painful than one could wish, and up to now only psql has really gotten it right among our supported client programs.

A new source-code subdirectory \texttt{src/fe\_utils/} has been created to hold this and other code that is shared across our client programs. Formerly such sharing was accomplished by symbolic linking or copying source files at build time, which was ugly and required duplicate compilation.

• Introduce \texttt{WaitEventSet} API to allow efficient waiting for event sets that usually do not change from one wait to the next (Andres Freund, Amit Kapila)

• Add a generic interface for writing WAL records (Alexander Korotkov, Petr Jelinek, Markus Nullmeier)
This change allows extensions to write WAL records for changes to pages using a standard layout. The problem of needing to replay WAL without access to the extension is solved by having generic replay code. This allows extensions to implement, for example, index access methods and have WAL support for them.

• Support generic WAL messages for logical decoding (Petr Jelínek, Andres Freund)

This feature allows extensions to insert data into the WAL stream that can be read by logical-decoding plugins, but is not connected to physical data restoration.

• Allow SP-GiST operator classes to store an arbitrary “traversal value” while descending the index (Alexander Lebedev, Teodor Sigaev)

This is somewhat like the “reconstructed value”, but it could be any arbitrary chunk of data, not necessarily of the same data type as the indexed column.

• Introduce a LOG_SERVER_ONLY message level for ereport() (David Steele)

This level acts like LOG except that the message is never sent to the client. It is meant for use in auditing and similar applications.

• Provide a Makefile target to build all generated headers (Michael Paquier, Tom Lane)

submake-generated-headers can now be invoked to ensure that generated backend header files are up-to-date. This is useful in subdirectories that might be built “standalone”.

• Support OpenSSL 1.1.0 (Andreas Karlsson, Heikki Linnakangas)

E.21.3.13. Additional Modules

• Add configuration parameter auto_explain.sample_rate to allow contrib/auto_explain to capture just a configurable fraction of all queries (Craig Ringer, Julien Rouhaud)

This allows reduction of overhead for heavy query traffic, while still getting useful information on average.

• Add contrib/bloom module that implements an index access method based on Bloom filtering (Teodor Sigaev, Alexander Korotkov)

This is primarily a proof-of-concept for non-core index access methods, but it could be useful in its own right for queries that search many columns.

• In contrib/cube, introduce distance operators for cubes, and support kNN-style searches in GiST indexes on cube columns (Stas Kelvich)

• Make contrib/hstore's hstore_to_jsonb_loose() and hstore_to_json_loose() functions agree on what is a number (Tom Lane)

Previously, hstore_to_jsonb_loose() would convert numeric-looking strings to JSON numbers, rather than strings, even if they did not exactly match the JSON syntax specification for numbers. This was inconsistent with hstore_to_json_loose(), so tighten the test to match the JSON syntax.

• Add selectivity estimation functions for contrib/intarray operators to improve plans for queries using those operators (Yury Zhuravlev, Alexander Korotkov)

• Make contrib/pageinspect's heap_page_items() function show the raw data in each tuple, and add new functions tuple_data_split() and heap_page_item_attrs() for inspection of individual tuple fields (Nikolay Shaplov)
• Add an optional S2K iteration count parameter to contrib/pgcrypto's pgp_sym_encrypt() function (Jeff Janes)

• Add support for “word similarity” to contrib/pg_trgm (Alexander Korotkov, Artur Zakirov)

  These functions and operators measure the similarity between one string and the most similar single word of another string.

• Add configuration parameter pg_trgm.similarity_threshold for contrib/pg_trgm's similarity threshold (Artur Zakirov)

  This threshold has always been configurable, but formerly it was controlled by special-purpose functions set_limit() and show_limit(). Those are now deprecated.

• Improve contrib/pg_trgm's GIN operator class to speed up index searches in which both common and rare keys appear (Jeff Janes)

• Improve performance of similarity searches in contrib/pg_trgm GIN indexes (Christophe Fornaroli)

• Add contrib/pg_visibility module to allow examining table visibility maps (Robert Haas)

• Add ssl_extension_info() function to contrib/sslinfo, to print information about SSL extensions present in the X509 certificate used for the current connection (Dmitry Voronin)

E.21.3.13.1. postgres_fdw

• Allow extension-provided operators and functions to be sent for remote execution, if the extension is whitelisted in the foreign server's options (Paul Ramsey)

  Users can enable this feature when the extension is known to exist in a compatible version in the remote database. It allows more efficient execution of queries involving extension operators.

• Consider performing sorts on the remote server (Ashutosh Bapat)

• Consider performing joins on the remote server (Shigeru Hanada, Ashutosh Bapat)

• When feasible, perform UPDATE or DELETE entirely on the remote server (Etsuro Fujita)

  Formerly, remote updates involved sending a SELECT FOR UPDATE command and then updating or deleting the selected rows one-by-one. While that is still necessary if the operation requires any local processing, it can now be done remotely if all elements of the query are safe to send to the remote server.

• Allow the fetch size to be set as a server or table option (Corey Huinker)

  Formerly, postgres_fdw always fetched 100 rows at a time from remote queries; now that behavior is configurable.

• Use a single foreign-server connection for local user IDs that all map to the same remote user (Ashutosh Bapat)

• Transmit query cancellation requests to the remote server (Michael Paquier, Etsuro Fujita)

  Previously, a local query cancellation request did not cause an already-sent remote query to terminate early.

E.22. Release 9.5.15

Release date: 2018-11-08
This release contains a variety of fixes from 9.5.14. For information about new features in the 9.5 major release, see Section E.37.

E.22.1. Migration to Version 9.5.15

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.13, see Section E.24.

E.22.2. Changes

• Fix corner-case failures in `has_foo_privilege()` family of functions (Tom Lane)
  
  Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. `has_column_privilege()` was additionally capable of crashing on some platforms.

• Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)

• Fix parsing of standard multi-character operators that are immediately followed by a comment or `+` or `-` (Andrew Gierth)
  
  This oversight could lead to parse errors, or to incorrect assignment of precedence.

• Avoid O(N^3) slowdown in lexer for long strings of `+` or `-` characters (Andrew Gierth)

• Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)

• Fix failure of `UPDATE/DELETE ... WHERE CURRENT OF ...` after rewinding the referenced cursor (Tom Lane)
  
  A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

• Fix `EvalPlanQual` to handle conditionally-executed InitPlans properly (Andrew Gierth, Tom Lane)
  
  This resulted in hard-to-reproduce crashes or wrong answers in concurrent updates, if they contained code such as an uncorrelated sub-`SELECT` inside a `CASE` construct.

• Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)
  
  This bug affected full-text-search operations, as well as `contrib/ltrie` and `contrib/pg_trgm`.

• Ensure that sequences owned by a foreign table are processed by `ALTER OWNER` on the table (Peter Eisentraut)
  
  The ownership change should propagate to such sequences as well, but this was missed for foreign tables.

• Ensure that the server will process already-received `NOTIFY` and `SIGTERM` interrupts before waiting for client input (Jeff Janes, Tom Lane)

• Fix over-allocation of space for `array_out()`’s result string (Keiichi Hirobe)

• Fix memory leak in repeated SP-GiST index scans (Tom Lane)
This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- Ensure that `ApplyLogicalMappingFile()` closes the mapping file when done with it (Tomas Vondra)

  Previously, the file descriptor was leaked, eventually resulting in failures during logical decoding.

- Fix logical decoding to handle cases where a mapped catalog table is repeatedly rewritten, e.g. by `VACUUM FULL` (Andres Freund)

- Prevent starting the server with `wal_level` set to too low a value to support an existing replication slot (Andres Freund)

- Avoid crash if a utility command causes infinite recursion (Tom Lane)

- When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master (Michael Paquier, Konstantin Knizhnik)

- Fix event triggers to handle nested `ALTER TABLE` commands (Michael Paquier, Álvaro Herrera)

- Propagate parent process's transaction and statement start timestamps to parallel workers (Konstantin Knizhnik)

  This prevents misbehavior of functions such as `transaction_timestamp()` when executed in a worker.

- Fix WAL file recycling logic to work correctly on standby servers (Michael Paquier)

  Depending on the setting of `archive_mode`, a standby might fail to remove some WAL files that could be removed.

- Fix handling of commit-timestamp tracking during recovery (Masahiko Sawasa, Michael Paquier)

  If commit timestamp tracking has been turned on or off, recovery might fail due to trying to fetch the commit timestamp for a transaction that did not record it.

- Randomize the `random()` seed in bootstrap and standalone backends, and in `initdb` (Noah Misch)

  The main practical effect of this change is that it avoids a scenario where `initdb` might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

- Allow DSM allocation to be interrupted (Chris Travers)

- Properly handle turning `full_page_writes` on dynamically (Kyotaro Horiguchi)

- Avoid possible buffer overrun when replaying GIN page recompression from WAL (Alexander Korotkov, Sivasubramanian Ramasubramanian)

- Fix missed fsync of a replication slot's directory (Konstantin Knizhnik, Michael Paquier)

- Fix unexpected timeouts when using `wal_sender_timeout` on a slow server (Noah Misch)

- Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)

  This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.

- Ensure background workers are stopped properly when the postmaster receives a fast-shutdown request before completing database startup (Alexander Kukushkin)
Release Notes

- Don’t run atexit callbacks when servicing SIGQUIT (Heikki Linnakangas)
- Don’t record foreign-server user mappings as members of extensions (Tom Lane)

If CREATE USER MAPPING is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can’t be extension members, so user mappings shouldn’t be either.

- Make syslogger more robust against failures in opening CSV log files (Tom Lane)
- Fix psql, as well as documentation examples, to call PQconsumeInput() before each PQnotifies() call (Tom Lane)

This fixes cases in which psql would not report receipt of a NOTIFY message until after the next command.

- Fix possible inconsistency in pg_dump’s sorting of dissimilar object names (Jacob Champion)
- Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)

This avoids failures due to the new policy of running restores with restrictive search path.

- Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)

pg_upgrade failed to preserve an event trigger’s extension-membership status.

- Fix pg_upgrade’s cluster state check to work correctly on a standby server (Bruce Momjian)
- Enforce type cube’s dimension limit in all contrib/cube functions (Andrey Borodin)

Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.

- Fix contrib/unaccent’s unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)

Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.

- Fix build problems on macOS 10.14 (Mojave) (Tom Lane)

Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.

It is now recommended that Perl-related extensions write $ perl_includespec rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.

Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.

- Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)

Perl no longer includes the current directory in its search path by default; work around that.

- On Windows, allow the regression tests to be run by an Administrator account (Andrew Dunstan)

To do this safely, pg_regress now gives up any such privileges at startup.

- Support building on Windows with Visual Studio 2015 or Visual Studio 2017 (Michael Paquier, Haribabu Kommi)
• Allow btree comparison functions to return INT_MIN (Tom Lane)

Up to now, we’ve forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of memcmp(), strcmp(), etc, as POSIX doesn’t place any such restriction on those functions. At least some recent versions of memcmp() can return INT_MIN, causing incorrect sort ordering. Hence, we’ve removed this restriction. Callers must now use the INVERT_COMPARE_RESULT() macro if they wish to invert the sort order.

• Fix recursion hazard in shared-invalidation message processing (Tom Lane)

This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.

This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.

• Save and restore SPI’s global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)

This prevents possible interference when one SPI-using function calls another.

• Provide ALLOCSET_DEFAULT_SIZES and sibling macros in back branches (Tom Lane)

These macros have existed since 9.6, but there were requests to add them to older branches to allow extensions to rely on them without branch-specific coding.

• Avoid using potentially-under-aligned page buffers (Tom Lane)

Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can’t place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.

• Make src/port/snprintf.c follow the C99 standard’s definition of snprintf()’s result value (Tom Lane)

On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.

• When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)

This avoids problems with missed floating point overflow checks.

• Fix configure’s detection of the result type of strerror_r() (Tom Lane)

The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.

• Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.23. Release 9.5.14

Release date: 2018-08-09

This release contains a variety of fixes from 9.5.13. For information about new features in the 9.5 major release, see Section E.37.
E.23.1. Migration to Version 9.5.14

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.13, see Section E.24.

E.23.2. Changes

• Fix failure to reset libpq's state fully between connection attempts (Tom Lane)

An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general, an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasichkov for reporting this issue. (CVE-2018-10915)

• Fix INSERT ... ON CONFLICT UPDATE through a view that isn't just SELECT * FROM ...

(Dean Rasheed, Amit Langote)

Erroneous expansion of an updatable view could lead to crashes or “attribute ... has the wrong type” errors, if the view's SELECT list doesn't match one-to-one with the underlying table's columns. Furthermore, this bug could be leveraged to allow updates of columns that an attacking user lacks UPDATE privilege for, if that user has INSERT and UPDATE privileges for some other column(s) of the table. Any user could also use it for disclosure of server memory. (CVE-2018-10925)

• Ensure that updates to the relfrozenxid and relminmxid values for “nailed” system catalogs are processed in a timely fashion (Andres Freund)

Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as pg_authid, because the stale cache data could persist into new sessions as well as existing ones.

• Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

• Avoid emitting a bogus WAL record when recycling an all-zero btree page (Amit Kapila)

This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

• During WAL replay, guard against corrupted record lengths exceeding 1GB (Michael Paquier)

Treat such a case as corrupt data. Previously, the code would try to allocate space and get a hard error, making recovery impossible.

• When ending recovery, delay writing the timeline history file as long as possible (Heikki Linnakangas)

This avoids some situations where a failure during recovery cleanup (such as a problem with a two-phase state file) led to inconsistent timeline state on-disk.

• Improve performance of WAL replay for transactions that drop many relations (Fujii Masao)
This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

• Improve performance of lock releasing in standby server WAL replay (Thomas Munro)

• Make logical WAL senders report streaming state correctly (Simon Riggs, Sawada Masahiko)

  The code previously mis-detected whether or not it had caught up with the upstream server.

• Fix bugs in snapshot handling during logical decoding, allowing wrong decoding results in rare cases (Arseny Sher, Álvaro Herrera)

• Ensure a table’s cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)

  Previously, the failed index’s OID could remain in the list, causing problems later in the same session.

• Fix mishandling of empty uncompressed posting list pages in GIN indexes (Sivasubramanian Ramasubramanian, Alexander Korotkov)

  This could result in an assertion failure after pg_upgrade of a pre-9.4 GIN index (9.4 and later will not create such pages).

• Ensure that VACUUM will respond to signals within btree page deletion loops (Andres Freund)

  Corrupted btree indexes could result in an infinite loop here, and that previously wasn’t interruptible without forcing a crash.

• Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)

  This resulted in failure to recognize that an index on a composite column could provide the sort order needed for a mergejoin on that column.

• Fix SQL-standard FETCH FIRST syntax to allow parameters ($n), as the standard expects (Andrew Gierth)

• Fix failure to schema-qualify some object names in getObjectDescription output (Kyotaro Horiguchi, Tom Lane)

  Names of collations, conversions, and text search objects were not schema-qualified when they should be.

• Widen COPY FROM’s current-line-number counter from 32 to 64 bits (David Rowley)

  This avoids two problems with input exceeding 4G lines: COPY FROM WITH HEADER would drop a line every 4G lines, not only the first line, and error reports could show a wrong line number.

• Add a string freeing function to ecpg’s pgtypes library, so that cross-module memory management problems can be avoided on Windows (Takayuki Tsunakawa)

  On Windows, crashes can ensue if the free call for a given chunk of memory is not made from the same DLL that malloc’ed the memory. The pgtypes library sometimes returns strings that it expects the caller to free, making it impossible to follow this rule. Add a PGTYPESchar_free() function that just wraps free, allowing applications to follow this rule.

• Fix ecpg’s support for long long variables on Windows, as well as other platforms that declare strtoll/strtoull nonstandardly or not at all (Dang Minh Huong, Tom Lane)

• Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in the semantics of a statement intra-session (Tom Lane)
This error led to assertion failures, or in rare cases, failure to enforce the INTO STRICT option as expected.

• Fix password prompting in client programs so that echo is properly disabled on Windows when stdin is not the terminal (Matthew Stickney)

• Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)

  The previous fix for quoting of search_path and other list-valued variables in pg_dump output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

• Fix pg_dump's failure to dump REPLICA IDENTITY properties for constraint indexes (Tom Lane)

  Manually created unique indexes were properly marked, but not those created by declaring UNIQUE or PRIMARY KEY constraints.

• Make pg_upgrade check that the old server was shut down cleanly (Bruce Momjian)

  The previous check could be fooled by an immediate-mode shutdown.

• Fix contrib/hstore_plperl to look through Perl scalar references, and to not crash if it doesn't find a hash reference where it expects one (Tom Lane)

• Fix crash in contrib/ltree's lca() function when the input array is empty (Pierre Ducroquet)

• Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

• Rearrange makefiles to ensure that programs link to freshly-built libraries (such as libpq.so) rather than ones that might exist in the system library directories (Tom Lane)

  This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

• Update time zone data files to tzdata release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.

  This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbreviations in use; the only known effect is that the is_dst column in the pg_timezone_names view will now be true in winter and false in summer in these cases.

E.24. Release 9.5.13

Release date: 2018-05-10

This release contains a variety of fixes from 9.5.12. For information about new features in the 9.5 major release, see Section E.37.

E.24.1. Migration to Version 9.5.13

A dump/restore is not required for those running 9.5.X.

However, if the function marking mistakes mentioned in the first changelog entry below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 9.5.12, see Section E.25.
E.24.2. Changes

• Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

The functions query_to_xml, cursor_to_xml, cursor_to_xmlschema, query_to_xmlschema, and query_to_xml_and_xmlschema should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions' pg_proc entries, for example ALTER FUNCTION pg_catalog.query_to_xml(text, boolean, boolean, text) VOLATILE. (Note that that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

• Avoid re-using TOAST value OIDs that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

Once the OID counter has wrapped around, it's possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value nnnnn” errors, which would persist until the dead entry was removed by VACUUM. Fix by not selecting such OIDs when creating a new TOAST entry.

• Change ANALYZE's algorithm for updating pg_class.reltuples (David Gould)

Previously, pages not actually scanned by ANALYZE were assumed to retain their old tuple density. In a large table where ANALYZE samples only a small fraction of the pages, this meant that the overall tuple density estimate could not change very much, so that reltuples would change nearly proportionally to changes in the table's physical size (retpages) regardless of what was actually happening in the table. This has been observed to result in reltuples becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that ANALYZE's sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

• Avoid deadlocks in concurrent CREATE INDEX CONCURRENTLY commands that are run under SERIALIZABLE or REPEATABLE READ transaction isolation (Tom Lane)

• Fix possible slow execution of REFRESH MATERIALIZED VIEW CONCURRENTLY (Thomas Munro)

• Fix UPDATE/DELETE ... WHERE CURRENT OF to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)

• Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)

This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.

• Fix possibly incorrect generation of an index-only-scan plan when the same table column appears in multiple index columns, and only some of those index columns use operator classes that can return the column value (Kyotaro Horiguchi)

• Fix misoptimization of CHECK constraints having provably-NULL subclauses of top-level AND/OR conditions (Tom Lane, Dean Rasheed)

This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.

• Fix executor crash due to double free in some GROUPING SET usages (Peter Geoghegan)
• Avoid crash if a table rewrite event trigger is added concurrently with a command that could call such a trigger (Álvaro Herrera, Andrew Gierth, Tom Lane)

• Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)

• Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)

• Fix overly strict sanity check in heap_prepare_freeze_tuple (Álvaro Herrera)

  This could result in incorrect “cannot freeze committed xmax” failures in databases that have been pg_upgrade'd from 9.2 or earlier.

• Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)

• Reduce locking during autovacuum worker scheduling (Jeff Janes)

  The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.

• Ensure client hostname is copied while copying pg_stat_activity data to local memory (Edmund Horner)

  Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the client hostname column to change unexpectedly if any existing session disconnected.

• Fix incorrect processing of multiple compound affixes in ispell dictionaries (Arthur Zakirov)

• Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on text columns (Tom Lane)

  Such searches would return the wrong set of rows in most non-C locales.

• Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)

  Previously, the tuple count was reported to be the same as that of the underlying table, which is wrong if the index is partial.

• Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)

  Previously it reported the estimated number of heap tuples, which might be inaccurate, and is certainly wrong if the index is partial.

• Fix a corner case where a streaming standby gets stuck at a WAL continuation record (Kyotaro Horiguchi)

• In logical decoding, avoid possible double processing of WAL data when a walsender restarts (Craig Ringer)

• Allow scalarltsel and scalargtsel to be used on non-core datatypes (Tomas Vondra)

• Reduce libpq's memory consumption when a server error is reported after a large amount of query output has been collected (Tom Lane)

  Discard the previous output before, not after, processing the error message. On some platforms, notably Linux, this can make a difference in the application's subsequent memory footprint.

• Fix double-free crashes in ecpg (Patrick Krecker, Jeevan Ladhe)

• Fix ecpg to handle long long int variables correctly in MSVC builds (Michael Meskes, Andrew Gierth)
• Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)

The local_preload_libraries, session_preload_libraries, shared_preload_libraries, and temp_tablespaces variables were not correctly quoted in pg_dump output. This would cause problems if settings for these variables appeared in CREATE FUNCTION ... SET or ALTER DATABASE/ROLE ... SET clauses.

• Fix pg_recvlogical to not fail against pre-v10 PostgreSQL servers (Michael Paquier)

A previous fix caused pg_recvlogical to issue a command regardless of server version, but it should only be issued to v10 and later servers.

• Ensure that pg_rewind deletes files on the target server if they are deleted from the source server during the run (Takayuki Tsunakawa)

Failure to do this could result in data inconsistency on the target, particularly if the file in question is a WAL segment.

• Fix pg_rewind to handle tables in non-default tablespaces correctly (Takayuki Tsunakawa)

• Fix overflow handling in PL/pgSQL integer FOR loops (Tom Lane)

The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

• Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisenstraut)

• Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)

• Support building with Microsoft Visual Studio 2015 (Michael Paquier)

Various fixes needed for VS2015 compatibility were previously back-patched into the 9.5 branch, but this one was missed.

• Rename internal b64_encode and b64_decode functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)

• Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)

This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.

• Update timezone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

E.25. Release 9.5.12

Release date: 2018-03-01

This release contains a variety of fixes from 9.5.11. For information about new features in the 9.5 major release, see Section E.37.

E.25.1. Migration to Version 9.5.12

A dump/restore is not required for those running 9.5.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that
you read the documentation changes described in the first changelog entry below, and take suitable steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 9.5.10, see Section E.27.

E.25.2. Changes

- Document how to configure installations and applications to guard against search-path-dependent trojan-horse attacks from other users (Noah Misch)

Using a `search_path` setting that includes any schemas writable by a hostile user enables that user to capture control of queries and then run arbitrary SQL code with the permissions of the attacked user. While it is possible to write queries that are proof against such hijacking, it is notationally tedious, and it's very easy to overlook holes. Therefore, we now recommend configurations in which no untrusted schemas appear in one's search path. Relevant documentation appears in Section 5.8.6 (for database administrators and users), Section 34.1 (for application authors), Section 38.16.1 (for extension authors), and `CREATE FUNCTION` (for authors of `SECURITY DEFINER` functions). (CVE-2018-1058)

- Avoid use of insecure `search_path` settings in `pg_dump` and other client programs (Noah Misch, Tom Lane)

`pg_dump`, `pg_upgrade`, `vacuumdb` and other PostgreSQL-provided applications were themselves vulnerable to the type of hijacking described in the previous changelog entry; since these applications are commonly run by superusers, they present particularly attractive targets. To make them secure whether or not the installation as a whole has been secured, modify them to include only the `pg_catalog` schema in their `search_path` settings. Autovacuum worker processes now do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example, user-provided functions in index expressions — the tighter `search_path` may result in errors, which will need to be corrected by adjusting those user-provided functions to not assume anything about what search path they are invoked under. That has always been good practice, but now it will be necessary for correct behavior. (CVE-2018-1058)

- Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom Lane)

If a CTE (`WITH` clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

- Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do not match mergeclauses” planner errors in corner cases.

- Repair `pg_upgrade`’s failure to preserve `relfrozenxid` for materialized views (Tom Lane, Andres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The problem would be more likely to occur in seldom-refreshed materialized views, or ones that were maintained only with `REFRESH MATERIALIZED VIEW CONCURRENTLY`.

If such corruption is observed, it can be repaired by refreshing the materialized view (without `CONCURRENTLY`).

2491
• Fix incorrect reporting of PL/Python function names in error CONTEXT stacks (Tom Lane)

An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function's name twice, rather than the expected results. Also, an error in a nested PL/Python DO block could result in a null pointer dereference crash on some platforms.

• Allow contrib/auto_explain's log_min_duration setting to range up to INT_MAX, or about 24 days instead of 35 minutes (Tom Lane)

E.26. Release 9.5.11

Release date: 2018-02-08

This release contains a variety of fixes from 9.5.10. For information about new features in the 9.5 major release, see Section E.37.

E.26.1. Migration to Version 9.5.11

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.10, see Section E.27.

E.26.2. Changes

• Ensure that all temporary files made by pg_upgrade are non-world-readable (Tom Lane, Noah Misch)

pg_upgrade normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing pg_dumpall -q output would be group- or world-readable, or even writable, if the user's umask setting allows. In typical usage on multi-user machines, the umask and/or the working directory's permissions would be tight enough to prevent problems; but there may be people using pg_upgrade in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

• Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

In some cases VACUUM would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

• Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)

These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.

• Fix incorrect query results from cases involving flattening of subqueries whose outputs are used in GROUPING SETS (Heikki Linnakangas)

• Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by ALTER TABLE NO INHERIT (Tom Lane)

• Fix spurious deadlock failures when multiple sessions are running CREATE INDEX CONCURRENTLY (Jeff Janes)

• Fix failures when an inheritance tree contains foreign child tables (Etsuro Fujita)

A mix of regular and foreign tables in an inheritance tree resulted in creation of incorrect plans for UPDATE and DELETE queries. This led to visible failures in some cases, notably when there are row-level triggers on a foreign child table.
• Repair failure with correlated sub-SELECT inside VALUES inside a LATERAL subquery (Tom Lane)

• Fix “could not devise a query plan for the given query” planner failure for some cases involving nested UNION ALL inside a lateral subquery (Tom Lane)

• Fix logical decoding to correctly clean up disk files for crashed transactions (Atsushi Torikoshi)

Logical decoding may spill WAL records to disk for transactions generating many WAL records. Normally these files are cleaned up after the transaction's commit or abort record arrives; but if no such record is ever seen, the removal code misbehaved.

• Fix walsender timeout failure and failure to respond to interrupts when processing a large transaction (Petr Jelinek)

• Fix has_sequence_privilege() to support WITH GRANT OPTION tests, as other privilege-testing functions do (Joe Conway)

• In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)

We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don't promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only xpath() and related functions; other XML code paths already acted this way.

• Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)

Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.

• Cope with failure to start a parallel worker process (Amit Kapila, Robert Haas)

Parallel query previously tended to hang indefinitely if a worker could not be started, as the result of fork() failure or other low-probability problems.

• Avoid unsafe alignment assumptions when working with __int128 (Tom Lane)

Typically, compilers assume that __int128 variables are aligned on 16-byte boundaries, but our memory allocation infrastructure isn't prepared to guarantee that, and increasing the setting of MAXALIGN seems infeasible for multiple reasons. Adjust the code to allow use of __int128 only when we can tell the compiler to assume lesser alignment. The only known symptom of this problem so far is crashes in some parallel aggregation queries.

• Prevent stack-overflow crashes when planning extremely deeply nested set operations (UNION/INTERSECT/EXCEPT) (Tom Lane)

• Fix null-pointer crashes for some types of LDAP URLs appearing in pg_hba.conf (Thomas Munro)

• Fix sample INSTR() functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)

These functions are stated to be Oracle® compatible, but they weren't exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.
The sample code has been adjusted to match Oracle's behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

- Fix `pg_dump` to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with `DATABASE` so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with `LARGE OBJECT`. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.

Note that this change has user-visible results in the output of `pg_restore --list`.

- Rename `pg_rewind`'s `copy_file_range` function to avoid conflict with new Linux system call of that name (Andres Freund)

This change prevents build failures with newer glibc versions.

- In `ecpg`, detect indicator arrays that do not have the correct length and report an error (David Rader)

- Avoid triggering a libc assertion in `contrib/hstore`, due to use of `memcpy()` with equal source and destination pointers (Tomas Vondra)

- Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

The scripts in `contrib/start-scripts/osx` use infrastructure that's been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory `contrib/start-scripts/macos` containing scripts that use the newer `launchd` infrastructure.

- Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dunstan)

- Support linking to MinGW-built versions of `libperl` (Noah Misch)

This allows building PL/Perl with some common Perl distributions for Windows.

- Fix MSVC build to test whether 32-bit `libperl` needs `-D_USE_32BIT_TIME_T` (Noah Misch)

Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.

- On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)

This may allow collection of a core dump for some early-startup failures that did not produce a dump before.

- On Windows, avoid encoding-conversion-related crashes when emitting messages very early in postmaster startup (Takayuki Tsunakawa)

- Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)

- Add support for spinlocks on Motorola 88K (David Carlier)

- Update time zone data files to `tzdata` release 2018c for DST law changes in Brazil, Sao Tome and Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The `US/Pacific-New` zone has been removed (it was only an alias for `America/Los_Angeles` anyway).
E.27. Release 9.5.10

Release date: 2017-11-09

This release contains a variety of fixes from 9.5.9. For information about new features in the 9.5 major release, see Section E.37.

E.27.1. Migration to Version 9.5.10

A dump/restore is not required for those running 9.5.X.

However, if you use BRIN indexes, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.5.8, see Section E.29.

E.27.2. Changes

• Ensure that INSERT ... ON CONFLICT DO UPDATE checks table permissions and RLS policies in all cases (Dean Rasheed)

  The update path of INSERT ... ON CONFLICT DO UPDATE requires SELECT permission on the columns of the arbiter index, but it failed to check for that in the case of an arbiter specified by constraint name. In addition, for a table with row level security enabled, it failed to check updated rows against the table's SELECT policies (regardless of how the arbiter index was specified). (CVE-2017-15099)

• Fix crash due to rowtype mismatch in json(b)_populate_recordset() (Michael Paquier, Tom Lane)

  These functions used the result rowtype specified in the FROM ... AS clause without checking that it matched the actual rowtype of the supplied tuple value. If it didn’t, that would usually result in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)

• Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)

  Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making $PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

  By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

• Fix BRIN index summarization to handle concurrent table extension correctly (Álvaro Herrera)

  Previously, a race condition allowed some table rows to be omitted from the index. It may be necessary to reindex existing BRIN indexes to recover from past occurrences of this problem.

• Fix possible failures during concurrent updates of a BRIN index (Tom Lane)

  These race conditions could result in errors like “invalid index offnum” or “inconsistent range map”.

• Fix crash when logical decoding is invoked from a SPI-using function, in particular any function written in a PL language (Tom Lane)

• Fix json_build_array(), json_build_object(), and their jsonb equivalents to handle explicit VARIADIC arguments correctly (Michael Paquier)

• Properly reject attempts to convert infinite float values to type numeric (Tom Lane, KaiGai Kohei)

2495
Previously the behavior was platform-dependent.

- Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)

- Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes (Tom Lane)

  Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

- Correctly detect hashability of range data types (Tom Lane)

  The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range's subtype has hash support. This does not affect any of the built-in range types, since they're all hashable anyway.

- Correctly ignore RelabelType expression nodes when determining relation distinctness (David Rowley)

  This allows the intended optimization to occur when a subquery has a result column of type varchar.

- Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)

  If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.

- Avoid SIGBUS crash on Linux when a DSM memory request exceeds the space available in tmpfs (Thomas Munro)

- Prevent low-probability crash in processing of nested trigger firings (Tom Lane)

- Allow COPY's FREEZE option to work when the transaction isolation level is REPEATABLE READ or higher (Noah Misch)

  This case was unintentionally broken by a previous bug fix.

- Correctly restore the umask setting when file creation fails in COPY or lo_export() (Peter Eisentraut)

- Give a better error message for duplicate column names in ANALYZE (Nathan Bossart)

- Fix mis-parsing of the last line in a non-newline-terminated pg_hba.conf file (Tom Lane)

- Fix pg_basebackup's matching of tablespace paths to canonicalize both paths before comparing (Michael Paquier)

  This is particularly helpful on Windows.

- Fix libpq to not require user's home directory to exist (Tom Lane)

  In v10, failure to find the home directory while trying to read ~/.pgpass was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading ~/.pg_service.conf, though this was less obvious since that file is not sought unless a service name is specified.

- Fix libpq to guard against integer overflow in the row count of a PGresult (Michael Paquier)

- Fix ecpg's handling of out-of-scope cursor declarations with pointer or array variables (Michael Meskes)
• In ecpglib, correctly handle backslashes in string literals depending on whether standard_conforming_strings is set (Tsunakawa Takayuki)

• Make ecpglib's Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengqi, Michael Meskes)

• Fix missing temp-install prerequisites for check-like Make targets (Noah Misch)

Some non-default test procedures that are meant to work like make check failed to ensure that the temporary installation was up to date.

• Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

This fixes various issues; the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no posixrules file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.

• Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

E.28. Release 9.5.9

Release date: 2017-08-31

This release contains a small number of fixes from 9.5.8. For information about new features in the 9.5 major release, see Section E.37.

E.28.1. Migration to Version 9.5.9

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.8, see Section E.29.

E.28.2. Changes

• Show foreign tables in information_schema.table_privileges view (Peter Eisentraut)

All other relevant information_schema views include foreign tables, but this one ignored them.

Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:

```sql
SET search_path TO information_schema;
CREATE OR REPLACE VIEW table_privileges AS
SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
       CAST(grantee.rolname AS sql_identifier) AS grantee,
       CAST(current_database() AS sql_identifier) AS table_catalog,
       CAST(nc.nspname AS sql_identifier) AS table_schema,
       CAST(c.relname AS sql_identifier) AS table_name,
       CAST(c.prtype AS character_data) AS privilege_type,
       CAST(CASE WHEN
                   -- object owner always has grant options
                   pg_has_role(grantee.oid, c.relowner, 'USAGE')
               OR c.grantable
```
Release Notes

THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
    CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE 'NO' END AS yes_or_no) AS with_hierarchy
FROM (SELECT oid, relname, relnamespace, relkind, relowner,
    (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class
) AS c (oid, relname, relnamespace, relkind, relowner,
    grantor, grantee, prtype, grantable),
pg_namespace nc,
pg_authid u_grantor,
( 
    SELECT oid, rolname FROM pg_authid
    UNION ALL
    SELECT 0::oid, 'PUBLIC'
) AS grantee (oid, rolname)
WHERE c.relnamespace = nc.oid
    AND c.relkind IN ('r', 'v', 'f')
    AND c.grantee = grantee.oid
    AND c.grantor = u_grantor.oid
    AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE',
        'DELETE', 'TRUNCATE', 'REFERENCES', 'TRIGGER')
    AND (pg_has_role(u_grantor.oid, 'USAGE')
        OR pg_has_role(grantee.oid, 'USAGE')
        OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

- Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)

  This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about "cannot drop active portal".

- Remove assertion that could trigger during a fatal exit (Tom Lane)

- Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

  Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant values that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.

- Fix crash in pg_restore when using parallel mode and using a list file to select a subset of items to restore (Fabrizio de Royes Mello)

- Change ecpg's parser to allow RETURNING clauses without attached C variables (Michael Meskes)

  This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

- Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

  This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that
Release Notes

_USE_32BIT_TIME_T is no longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

- Fix make check to behave correctly when invoked via a non-GNU make program (Thomas Munro)

E.29. Release 9.5.8

Release date: 2017-08-10

This release contains a variety of fixes from 9.5.7. For information about new features in the 9.5 major release, see Section E.37.

E.29.1. Migration to Version 9.5.8

A dump/restore is not required for those running 9.5.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.5.7, see Section E.30.

E.29.2. Changes

- Further restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Noah Misch)

The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have USAGE permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since information_schema.user_mapping_options does not show the options in such cases, pg_user_mappings should not either. (CVE-2017-7547)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, you will need to do the following:

1. Restart the postmaster after adding allow_system_table_mods = true to postgresql.conf. (In versions supporting ALTER SYSTEM, you can use that to make the configuration change, but you'll still need a restart.)

2. In each database of the cluster, run the following commands as superuser:

```sql
SET search_path = pg_catalog;
CREATE OR REPLACE VIEW pg_user_mappings AS
SELECT
  U.oid       AS umid,
  S.oid       AS srvid,
  S.srvname   AS srvname,
  U.umuser    AS umuser,
  CASE WHEN U.umuser = 0 THEN 'public'
    ELSE A.rolname
  END AS usename,
FROM pg_user_mappings U
JOIN pg_user_property_optum S
ON U.umpropertyopt = S.oid
JOIN pg_rolename A
ON A.oid = U.rolname
WHERE U.umuser = 0 OR A.rolname = current_user
AND (pg_has_role(S.srvowner, 'USAGE')
  OR has_server_privilege(S.oid,
    'USAGE'))
```

2499
Release Notes

OR (U.umuser = 0 AND pg_has_role(S.srvowner, 'USAGE'))
OR (SELECT rolsuper FROM pg_authid WHERE rolname = current_user)
THEN U.umoptions
ELSE NULL END AS umoptions
FROM pg_user_mapping U
LEFT JOIN pg_authid A ON (A.oid = U.umuser) JOIN pg_foreign_server S ON (U.umserver = S.oid);

3. Do not forget to include the template0 and template1 databases, or the vulnerability will still exist in databases you create later. To fix template0, you'll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;
and then after fixing template0, undo that with

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;
In prior versions, instead use

UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';

4. Finally, remove the allow_system_table_mods configuration setting, and again restart the postmaster.

- Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)
libpq ignores empty password specifications, and does not transmit them to the server. So, if a user's password has been set to the empty string, it's impossible to log in with that password via psql or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, md5, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

- Make lo_put() check for UPDATE privilege on the target large object (Tom Lane, Michael Paquier)
lo_put() should surely require the same permissions as lowrite(), but the check was missing, allowing any user to change the data in a large object. (CVE-2017-7548)

- Correct the documentation about the process for upgrading standby servers with pg_upgrade (Bruce Momjian)
The previous documentation instructed users to start/stop the primary server after running pg_upgrade but before syncing the standby servers. This sequence is unsafe.

- Fix concurrent locking of tuple update chains (Álvaro Herrera)
If several sessions concurrently lock a tuple update chain with nonconflicting lock modes using an old snapshot, and they all succeed, it was possible for some of them to nonetheless fail (and conclude there is no live tuple version) due to a race condition. This had consequences such as foreign-key checks failing to see a tuple that definitely exists but is being updated concurrently.
• Fix potential data corruption when freezing a tuple whose XMAX is a multixact with exactly one
  still-interesting member (Teodor Sigaev)

• Avoid integer overflow and ensuing crash when sorting more than one billion tuples in-memory
  (Sergey Koposov)

• On Windows, retry process creation if we fail to reserve the address range for our shared memory
  in the new process (Tom Lane, Amit Kapila)
  
  This is expected to fix infrequent child-process-launch failures that are probably due to interference
  from antivirus products.

• Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas
  Munro, Tom Lane)

• Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael
  Paquier)

• Prevent sending SSL session tickets to clients (Tom Lane)
  
  This fix prevents reconnection failures with ticket-aware client-side SSL code.

• Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)

• Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and im-
  mediate restart (Tom Lane)
  
  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively
  ignored.

• Ensure that the statistics collector's receive buffer size is at least 100KB (Tom Lane)
  
  This reduces the risk of dropped statistics data on older platforms whose default receive buffer size
  is less than that.

• Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes
  an XLOG_SWITCH WAL record (Andres Freund)

• Fix walsender to exit promptly when client requests shutdown (Tom Lane)

• Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)

• Prevent walsender-triggered panics during shutdown checkpoints (Andres Freund, Michael
  Paquier)

• Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom
  Lane)

• Fix leakage of small subtransactions spilled to disk during logical decoding (Andres Freund)
  
  This resulted in temporary files consuming excessive disk space.

• Reduce the work needed to build snapshots during creation of logical-decoding slots (Andres Fre-
  und, Petr Jelinek)
  
  The previous algorithm was infeasibly expensive on a server with a lot of open transactions.

• Fix race condition that could indefinitely delay creation of logical-decoding slots (Andres Freund,
  Petr Jelinek)

• Reduce overhead in processing syscache invalidation events (Tom Lane)
This is particularly helpful for logical decoding, which triggers frequent cache invalidation.

- Fix cases where an `INSERT` or `UPDATE` assigns to more than one element of a column that is of domain-over-array type (Tom Lane)

- Allow window functions to be used in sub-`SELECT`s that are within the arguments of an aggregate function (Tom Lane)

- Move autogenerated array types out of the way during `ALTER ... RENAME` (Vik Fearing)

  Previously, we would rename a conflicting autogenerated array type out of the way during `CREATE`; this fix extends that behavior to renaming operations.

- Fix dangling pointer in `ALTER TABLE` when there is a comment on a constraint belonging to the table (David Rowley)

  Re-applying the comment to the reconstructed constraint could fail with a weird error message, or even crash.

- Ensure that `ALTER USER ... SET` accepts all the syntax variants that `ALTER ROLE ... SET` does (Peter Eisentraut)

- Properly update dependency info when changing a datatype I/O function's argument or return type from `opaque` to the correct type (Heikki Linnakangas)

  `CREATE TYPE` updates I/O functions declared in this long-obsolete style, but it forgot to record a dependency on the type, allowing a subsequent `DROP TYPE` to leave broken function definitions behind.

- Reduce memory usage when `ANALYZE` processes a `tsvector` column (Heikki Linnakangas)

- Fix unnecessary precision loss and sloppy rounding when multiplying or dividing `money` values by integers or floats (Tom Lane)

- Tighten checks for whitespace in functions that parse identifiers, such as `regprocedurein()` (Tom Lane)

  Depending on the prevailing locale, these functions could misinterpret fragments of multibyte characters as whitespace.

- Use relevant `#define` symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)

  This avoids portability problems, typically manifesting as a “handshake” mismatch during library load, when working with recent Perl versions.

- In `libpq`, reset GSS/SASL and SSPI authentication state properly after a failed connection attempt (Michael Paquier)

  Failure to do this meant that when falling back from SSL to non-SSL connections, a GSS/SASL failure in the SSL attempt would always cause the non-SSL attempt to fail. SSPI did not fail, but it leaked memory.

- In `psql`, fix failure when `COPY FROM STDIN` is ended with a keyboard EOF signal and then another `COPY FROM STDIN` is attempted (Thomas Munro)

  This misbehavior was observed on BSD-derived platforms (including macOS), but not on most others.

- Fix `pg_dump` and `pg_restore` to emit `REFRESH MATERIALIZED VIEW` commands last (Tom Lane)
Release Notes

This prevents errors during dump/restore when a materialized view refers to tables owned by a different user.

- Improve pg_dump/pg_restore's reporting of error conditions originating in zlib (Vladimir Kun-schikov, Álvaro Herrera)

- Fix pg_dump with the --clean option to drop event triggers as expected (Tom Lane)
  
  It also now correctly assigns ownership of event triggers; before, they were restored as being owned by the superuser running the restore script.

- Fix pg_dump to not emit invalid SQL for an empty operator class (Daniel Gustafsson)

- Fix pg_dump output to stdout on Windows (Kuntal Ghosh)
  
  A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_get_ruledef() to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)
  
  In some corner cases, pg_dump relies on pg_get_ruledef() to dump views, so that this error could result in dump/reload failures.

- Fix dumping of outer joins with empty constraints, such as the result of a NATURAL LEFT JOIN with no common columns (Tom Lane)

- Fix dumping of function expressions in the FROM clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)

- Fix pg_basebackup output to stdout on Windows (Haribabu Kommi)
  
  A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_rewind to correctly handle files exceeding 2GB (Kuntal Ghosh, Michael Paquier)
  
  Ordinarily such files won't appear in PostgreSQL data directories, but they could be present in some cases.

- Fix pg_upgrade to ensure that the ending WAL record does not have walLevel = minimum (Bruce Momjian)
  
  This condition could prevent upgraded standby servers from reconnecting.

- Fix pg_xlogdump's computation of WAL record length (Andres Freund)

- In postgres_fdw, re-establish connections to remote servers after ALTER SERVER or ALTER USER MAPPING commands (Kyotaro Horiguchi)
  
  This ensures that option changes affecting connection parameters will be applied promptly.

- In postgres_fdw, allow cancellation of remote transaction control commands (Robert Haas, Rafia Sabih)
  
  This change allows us to quickly escape a wait for an unresponsive remote server in many more cases than previously.

- Increase MAX_SYSCACHE_CALLBACKS to provide more room for extensions (Tom Lane)

- Always use -fPIC, not -fpic, when building shared libraries with gcc (Tom Lane)
This supports larger extension libraries on platforms where it makes a difference.

- Fix unescaped-braces issue in our build scripts for Microsoft MSVC, to avoid a warning or error from recent Perl versions (Andrew Dunstan)

- In MSVC builds, handle the case where the OpenSSL library is not within a VC subdirectory (Andrew Dunstan)

- In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)

  This fixes a former need to move things around in standard Windows installations of libxml2.

- In MSVC builds, recognize a Tcl library that is named tcl86.lib (Noah Misch)

- In MSVC builds, honor PROVE_FLAGS settings on vcregress.pl's command line (Andrew Dunstan)

E.30. Release 9.5.7

Release date: 2017-05-11

This release contains a variety of fixes from 9.5.6. For information about new features in the 9.5 major release, see Section E.37.

E.30.1. Migration to Version 9.5.7

A dump/restore is not required for those running 9.5.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are using third-party replication tools that depend on “logical decoding”, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.5.6, see Section E.31.

E.30.2. Changes

- Restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)

  The previous coding allowed the owner of a foreign server object, or anyone he has granted server USAGE permission to, to see the options for all user mappings associated with that server. This might well include passwords for other users. Adjust the view definition to match the behavior of information_schema.user_mapping_options, namely that these options are visible to the user being mapped, or if the mapping is for PUBLIC and the current user is the server owner, or if the current user is a superuser. (CVE-2017-7486)

  By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, follow the corrected procedure shown in the changelog entry for CVE-2017-7547, in Section E.29.

- Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

  Some selectivity estimation functions in the planner will apply user-defined operators to values obtained from pg_statistic, such as most common values and histogram entries. This occurs before table permissions are checked, so a nefarious user could exploit the behavior to obtain these
values for table columns he does not have permission to read. To fix, fall back to a default estimate if the operator's implementation function is not certified leak-proof and the calling user does not have permission to read the table column whose statistics are needed. At least one of these criteria is satisfied in most cases in practice. (CVE-2017-7484)

• Restore libpq's recognition of the PGREQUIRESSL environment variable (Daniel Gustafsson)

Processing of this environment variable was unintentionally dropped in PostgreSQL 9.3, but its documentation remained. This creates a security hazard, since users might be relying on the environment variable to force SSL-encrypted connections, but that would no longer be guaranteed. Restore handling of the variable, but give it lower priority than PGSSLMODE, to avoid breaking configurations that work correctly with post-9.3 code. (CVE-2017-7485)

• Fix possibly-invalid initial snapshot during logical decoding (Petr Jelinek, Andres Freund)

The initial snapshot created for a logical decoding replication slot was potentially incorrect. This could cause third-party tools that use logical decoding to copy incomplete/inconsistent initial data. This was more likely to happen if the source server was busy at the time of slot creation, or if another logical slot already existed.

If you are using a replication tool that depends on logical decoding, and it should have copied a nonempty data set at the start of replication, it is advisable to recreate the replica after installing this update, or to verify its contents against the source server.

• Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)

This could result in an unlogged index being set to an invalid state after a crash and restart. Such a problem would persist until the index was dropped and rebuilt.

• Fix incorrect reconstruction of pg_subtrans entries when a standby server replays a prepared but uncommitted two-phase transaction (Tom Lane)

In most cases this turned out to have no visible ill effects, but in corner cases it could result in circular references in pg_subtrans, potentially causing infinite loops in queries that examine rows modified by the two-phase transaction.

• Avoid possible crash in walsender due to failure to initialize a string buffer (Stas Kelvich, Fujii Masao)

• Fix possible crash when rescanning a nearest-neighbor index-only scan on a GiST index (Tom Lane)

• Fix postmaster's handling of fork() failure for a background worker process (Tom Lane)

Previously, the postmaster updated portions of its state as though the process had been launched successfully, resulting in subsequent confusion.

• Fix crash or wrong answers when a GROUPING SETS column's data type is hashable but not sortable (Pavan Deolasee)

• Avoid applying “physical targetlist” optimization to custom scans (Dmitry Ivanov, Tom Lane)

This optimization supposed that retrieving all columns of a tuple is inexpensive, which is true for ordinary Postgres tuples; but it might not be the case for a custom scan provider.

• Use the correct sub-expression when applying a FOR ALL row-level-security policy (Stephen Frost)

In some cases the WITH CHECK restriction would be applied when the USING restriction is more appropriate.

• Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL (Julien Rouhaud, Tom Lane)
Due to lack of a cache flush step between commands in an extension script file, non-utility queries might not see the effects of an immediately preceding catalog change, such as `ALTER TABLE ... RENAME`.

- **Skip tablespace privilege checks when `ALTER TABLE ... ALTER COLUMN TYPE` rebuilds an existing index (Noah Misch)**

  The command failed if the calling user did not currently have `CREATE` privilege for the tablespace containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be rebuilt where it is.

- **Fix `ALTER TABLE ... VALIDATE CONSTRAINT` to not recurse to child tables when the constraint is marked `NO INHERIT` (Amit Langote)**

  This fix prevents unwanted "constraint does not exist" failures when no matching constraint is present in the child tables.

- **Avoid dangling pointer in `COPY ... TO` when row-level security is active for the source table (Tom Lane)**

  Usually this had no ill effects, but sometimes it would cause unexpected errors or crashes.

- **Avoid accessing an already-closed relcache entry in `CLUSTER` and `VACUUM FULL` (Tom Lane)**

  With some bad luck, this could lead to indexes on the target relation getting rebuilt with the wrong persistence setting.

- **Fix `VACUUM` to account properly for pages that could not be scanned due to conflicting page pins (Andrew Gierth)**

  This tended to lead to underestimation of the number of tuples in the table. In the worst case of a small heavily-contended table, `VACUUM` could incorrectly report that the table contained no tuples, leading to very bad planning choices.

- **Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests (Tom Lane, Thomas Munro)**

- **Fix integer-overflow problems in `interval` comparison (Kyotaro Horiguchi, Tom Lane)**

  The comparison operators for type `interval` could yield wrong answers for intervals larger than about 296000 years. Indexes on columns containing such large values should be reindexed, since they may be corrupt.

- **Fix `cursor_to_xml()` to produce valid output with `tableforest = false` (Thomas Munro, Peter Eisentraut)**

  Previously it failed to produce a wrapping `<table>` element.

- **Fix roundoff problems in `float8_timestamptz()` and `make_interval()` (Tom Lane)**

  These functions truncated, rather than rounded, when converting a floating-point value to integer microseconds; that could cause unexpectedly off-by-one results.

- **Fix `pg_get_object_address()` to handle members of operator families correctly (Álvaro Herrera)**

- **Improve performance of `pg_timezone_names` view (Tom Lane, David Rowley)**

- **Reduce memory management overhead for contexts containing many large blocks (Tom Lane)**

- **Fix sloppy handling of corner-case errors from `lseek()` and `close()` (Tom Lane)**
Neither of these system calls are likely to fail in typical situations, but if they did, `fd.c` could get quite confused.

- Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)
  
  This could result in attempting to write to the event log when that isn’t accessible, so that no logging happens at all.

- Fix `ecpg` to support `COMMIT PREPARED` and `ROLLBACK PREPARED` (Masahiko Sawada)

- Fix a double-free error when processing dollar-quoted string literals in `ecpg` (Michael Meskes)

- In `pg_dump`, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)
  
  In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.

- Avoid emitting an invalid list file in `pg_restore -l` when SQL object names contain newlines (Tom Lane)
  
  Replace newlines by spaces, which is sufficient to make the output valid for `pg_restore -L`'s purposes.

- Fix `pg_upgrade` to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)
  
  Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.

- Improve error handling in `contrib/adminpack's pg_file_write()` function (Noah Misch)
  
  Notably, it failed to detect errors reported by `fclose()`.

- In `contrib/dblink`, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

- Fix `contrib/pg_trgm`'s extraction of trigrams from regular expressions (Tom Lane)
  
  In some cases it would produce a broken data structure that could never match anything, leading to GIN or GiST indexes scans that use a trigram index not finding any matches to the regular expression.

- In `contrib/postgres_fdw`, transmit query cancellation requests to the remote server (Michael Paquier, Etsuro Fujita)
  
  Previously, a local query cancellation request did not cause an already-sent remote query to terminate early. This is a back-patch of work originally done for 9.6.

- Support Tcl 8.6 in MSVC builds (Álvaro Herrera)

- Sync our copy of the timezone library with IANA release tzcode2017b (Tom Lane)
  
  This fixes a bug affecting some DST transitions in January 2038.

- Update time zone data files to tzdata release 2017b for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.
The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the `pg_timezone_names` view nor used for output.

- Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds (David Rowley)

The Microsoft MSVC build scripts neglected to install the `posixrules` file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the US/Eastern zone. If you don't want that, remove the `posixrules` file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.

### E.31. Release 9.5.6

**Release date:** 2017-02-09

This release contains a variety of fixes from 9.5.5. For information about new features in the 9.5 major release, see Section E.37.

#### E.31.1. Migration to Version 9.5.6

A dump/restore is not required for those running 9.5.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.5.5, see Section E.32.

#### E.31.2. Changes

- Fix a race condition that could cause indexes built with `CREATE INDEX CONCURRENTLY` to be corrupt (Pavan Deolasee, Tom Lane)

  If `CREATE INDEX CONCURRENTLY` was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the `CREATE INDEX` command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

- Ensure that the special snapshot used for catalog scans is not invalidated by premature data pruning (Tom Lane)

  Backends failed to account for this snapshot when advertising their oldest xmin, potentially allowing concurrent vacuuming operations to remove data that was still needed. This led to transient failures along the lines of “cache lookup failed for relation 1255”.

- Fix incorrect WAL logging for BRIN indexes (Kuntal Ghosh)

  The WAL record emitted for a BRIN “revmap” page when moving an index tuple to a different page was incorrect. Replay would make the related portion of the index useless, forcing it to be recomputed.

- Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)
Previously, this was skipped when wal_level = minimal, but actually it's necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

- Reduce interlocking on standby servers during the replay of btree index vacuuming operations (Simon Riggs)

  This change avoids substantial replication delays that sometimes occurred while replaying such operations.

- If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)

- Ensure that hot standby feedback works correctly when it's enabled at standby server start (Ants Aasma, Craig Ringer)

- Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

- Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)

  This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

- Fix check for when an extension member object can be dropped (Tom Lane)

  Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

- Make sure ALTER TABLE preserves index tablespace assignments when rebuilding indexes (Tom Lane, Michael Paquier)

  Previously, non-default settings of default_tablespace could result in broken indexes.

- Fix incorrect updating of trigger function properties when changing a foreign-key constraint's deferrability properties with ALTER TABLE ... ALTER CONSTRAINT (Tom Lane)

  This led to odd failures during subsequent exercise of the foreign key, as the triggers were fired at the wrong times.

- Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)

  This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.

- Fix ALTER TABLE ... SET DATA TYPE ... USING when child table has different column ordering than the parent (Álvaro Herrera)

  Failure to adjust the column numbering in the USING expression led to errors, typically “attribute N has wrong type”.

- Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via ALTER TABLE ... INHERIT (Amit Langote)

  The OID column should be treated the same as regular user columns in this case, but it wasn't, leading to odd behavior in later inheritance changes.

- Fix CREATE OR REPLACE VIEW to update the view query before attempting to apply the new view options (Dean Rasheed)

  Previously the command would fail if the new options were inconsistent with the old view definition.

- Report correct object identity during ALTER TEXT SEARCH CONFIGURATION (Artur Zakirov)

  The wrong catalog OID was reported to extensions such as logical decoding.
• Fix commit timestamp mechanism to not fail when queried about the special XIDs Frozen-TransactionId and BootstrapTransactionId (Craig Ringer)

• Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

• Fix incorrect use of view reloptions as regular table reloptions (Tom Lane)

The symptom was spurious “ON CONFLICT is not supported on table ... used as a catalog table” errors when the target of INSERT ... ON CONFLICT is a view with cascade option.

• Fix incorrect “target lists can have at most N entries” complaint when using ON CONFLICT with wide tables (Tom Lane)

• Prevent multicolumn expansion of foo.* in an UPDATE source expression (Tom Lane)

This led to “UPDATE target count mismatch --- internal error”. Now the syntax is understood as a whole-row variable, as it would be in other contexts.

• Ensure that column typmods are determined accurately for multi-row VALUES constructs (Tom Lane)

This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a varchar value) but later values don’t share the same limit.

• Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (U&'...') or Unicode identifier (U&"...").

• Ensure that a purely negative text search query, such as !foo, matches empty tsvector (Tom Dunstan)

Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

• Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

• Fix performance problems in ts_rewrite() (Tom Lane)

• Fix ts_rewrite()'s handling of nested NOT operators (Tom Lane)

• Improve speed of user-defined aggregates that use array_append() as transition function (Tom Lane)

• Fix array_fill() to handle empty arrays properly (Tom Lane)

• Fix possible crash in array_position() or array_positions() when processing arrays of records (Junseok Yang)

• Fix one-byte buffer overrun in quote_literal_cstr() (Heikki Linnakangas)

The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.
Release Notes

- Prevent multiple calls of `pg_start_backup()` and `pg_stop_backup()` from running concurrently (Michael Paquier)
  
  This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

- Disable transform that attempted to remove no-op `AT TIME ZONE` conversions (Tom Lane)
  
  This resulted in wrong answers when the simplified expression was used in an index condition.

- Avoid discarding `interval-to-interval` casts that aren't really no-ops (Tom Lane)

  In some cases, a cast that should result in zeroing out low-order `interval` fields was mistakenly deemed to be a no-op and discarded. An example is that casting from `INTERVAL MONTH` to `INTERVAL YEAR` failed to clear the months field.

- Fix bugs in transmitting GUC parameter values to parallel workers (Michael Paquier, Tom Lane)

- Ensure that cached plans are invalidated by changes in foreign-table options (Amit Langote, Etsuro Fujita, Ashutosh Bapat)

- Fix `pg_dump` to dump user-defined casts and transforms that use built-in functions (Stephen Frost)

- Fix `pg_restore` with `--create --if-exists` to behave more sanely if an archive contains unrecognized `DROP` commands (Tom Lane)

  This doesn't fix any live bug, but it may improve the behavior in future if `pg_restore` is used with an archive generated by a later `pg_dump` version.

- Fix `pg_basebackup`'s rate limiting in the presence of slow I/O (Antonin Houska)

  If disk I/O was transiently much slower than the specified rate limit, the calculation overflowed, effectively disabling the rate limit for the rest of the run.

- Fix `pg_basebackup`'s handling of symlinked `pg_stat_tmp` and `pg_replslot` subdirectories (Magnus Hagander, Michael Paquier)

- Fix possible `pg_basebackup` failure on standby server when including WAL files (Amit Kapila, Robert Haas)

- Fix possible mishandling of expanded arrays in domain check constraints and `CASE` execution (Tom Lane)

  It was possible for a PL/pgSQL function invoked in these contexts to modify or even delete an array value that needs to be preserved for additional operations.

- Fix nested uses of PL/pgSQL functions in contexts such as domain check constraints evaluated during assignment to a PL/pgSQL variable (Tom Lane)

- Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)

  This avoids failures if the objects are used after a Python garbage collection cycle has occurred.

- Fix PL/Tcl to support triggers on tables that have `.tupno` as a column name (Tom Lane)

  This matches the (previously undocumented) behavior of PL/Tcl's `spi_exec` and `spi_execp` commands, namely that a magic `.tupno` column is inserted only if there isn't a real column named that.

- Allow DOS-style line endings in `~/.pgpass` files, even on Unix (Vik Fearing)

  This change simplifies use of the same password file across Unix and Windows machines.
• Fix one-byte buffer overrun if ecpg is given a file name that ends with a dot (Takayuki Tsunakawa)
• Fix psql's tab completion for ALTER DEFAULT PRIVILEGES (Gilles Darold, Stephen Frost)
• In psql, treat an empty or all-blank setting of the PAGER environment variable as meaning “no pager” (Tom Lane)

Previously, such a setting caused output intended for the pager to vanish entirely.

• Improve contrib/dblink's reporting of low-level libpq errors, such as out-of-memory (Joe Conway)
• Teach contrib/dblink to ignore irrelevant server options when it uses a contrib/postgres_fdw foreign server as the source of connection options (Corey Huinker)

Previously, if the foreign server object had options that were not also libpq connection options, an error occurred.

• Fix portability problems in contrib/pageinspect's functions for GIN indexes (Peter Eisentraut, Tom Lane)
• On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)
• Sync our copy of the timezone library with IANA release tzcode2016j (Tom Lane)

This fixes various issues, most notably that timezone data installation failed if the target directory didn't support hard links.

• Update time zone data files to tzdata release 2016j for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarctica/Cassey. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.

E.32. Release 9.5.5

Release date: 2016-10-27

This release contains a variety of fixes from 9.5.4. For information about new features in the 9.5 major release, see Section E.37.

E.32.1. Migration to Version 9.5.5

A dump/restore is not required for those running 9.5.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted free space maps.

Also, if you are upgrading from a version earlier than 9.5.2, see Section E.35.

E.32.2. Changes

• Fix WAL-logging of truncation of relation free space maps and visibility maps (Pavan Deolasee, Heikki Linnakangas)

It was possible for these files to not be correctly restored during crash recovery, or to be written incorrectly on a standby server. Bogus entries in a free space map could lead to attempts to access pages that have been truncated away from the relation itself, typically producing errors like “could not read block XXX: read only 0 of 8192 bytes”. Checksum failures in the visibility map are also possible, if checksumming is enabled.
Procedures for determining whether there is a problem and repairing it if so are discussed at https://wiki.postgresql.org/wiki/Free_Space_Map_Problems.

- Fix incorrect creation of GIN index WAL records on big-endian machines (Tom Lane)
  The typical symptom was “unexpected GIN leaf action” errors during WAL replay.

- Fix SELECT FOR UPDATE/SHARE to correctly lock tuples that have been updated by a subsequently-aborted transaction (Álvaro Herrera)
  In 9.5 and later, the SELECT would sometimes fail to return such tuples at all. A failure has not been proven to occur in earlier releases, but might be possible with concurrent updates.

- Fix EvalPlanQual rechecks involving CTE scans (Tom Lane)
  The recheck would always see the CTE as returning no rows, typically leading to failure to update rows that were recently updated.

- Fix deletion of speculatively inserted TOAST tuples when backing out of INSERT ... ON CONFLICT (Oskari Saarenmaa)
  In the race condition where two transactions try to insert conflicting tuples at about the same time, the loser would fail with an “attempted to delete invisible tuple” error if its insertion included any TOAST'ed fields.

- Don’t throw serialization errors for self-conflicting insertions in INSERT ... ON CONFLICT (Thomas Munro, Peter Geoghegan)

- Fix improper repetition of previous results from hashed aggregation in a subquery (Andrew Gierth)
  The test to see if we can reuse a previously-computed hash table of the aggregate state values neglected the possibility of an outer query reference appearing in an aggregate argument expression. A change in the value of such a reference should lead to recalculating the hash table, but did not.

- Fix query-lifespan memory leak in a bulk UPDATE on a table with a PRIMARY KEY or REPLICA IDENTITY index (Tom Lane)

- Fix COPY with a column name list from a table that has row-level security enabled (Adam Brightwell)

- Fix EXPLAIN to emit valid XML when track_io_timing is on (Markus Winand)
  Previously the XML output-format option produced syntactically invalid tags such as <I/O-Read-Time>. That is now rendered as <I-O-Read-Time>.

- Suppress printing of zeroes for unmeasured times in EXPLAIN (Maksim Milyutin)
  Certain option combinations resulted in printing zero values for times that actually aren’t ever measured in that combination. Our general policy in EXPLAIN is not to print such fields at all, so do that consistently in all cases.

- Fix statistics update for TRUNCATE in a prepared transaction (Stas Kelvich)

- Fix timeout length when VACUUM is waiting for exclusive table lock so that it can truncate the table (Simon Riggs)
  The timeout was meant to be 50 milliseconds, but it was actually only 50 microseconds, causing VACUUM to give up on truncation much more easily than intended. Set it to the intended value.

- Fix bugs in merging inherited CHECK constraints while creating or altering a table (Tom Lane, Amit Langote)
Allow identical CHECK constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a NOT VALID constraint on the child. Likewise, prevent merging of a NO INHERIT child constraint with an inherited constraint.

- Show a sensible value in pg_settings.unit for min_wal_size and max_wal_size (Tom Lane)
- Remove artificial restrictions on the values accepted by numeric_in() and numeric_recv() (Tom Lane)
  
  We allow numeric values up to the limit of the storage format (more than 1e100000), so it seems fairly pointless that numeric_in() rejected scientific-notation exponents above 1000. Likewise, it was silly for numeric_recv() to reject more than 1000 digits in an input value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)
- Preserve commit timestamps across server restart (Julien Rouhaud, Craig Ringer)
  
  With track_commit_timestamp turned on, old commit timestamps became inaccessible after a clean server restart.

- Fix logical WAL decoding to work properly when a subtransaction's WAL output is large enough to spill to disk (Andres Freund)
- Fix possible sorting error when aborting use of abbreviated keys (Peter Geoghegan)
  
  In the worst case, this could result in a corrupt btree index, which would need to be rebuilt using REINDEX. However, the situation is believed to be rare.

- Fix file descriptor leakage when truncating a temporary relation of more than 1GB (Andres Freund)
- Disallow starting a standalone backend with standby_mode turned on (Michael Paquier)
  
  This can't do anything useful, since there will be no WAL receiver process to fetch more WAL data; and it could result in misbehavior in code that wasn't designed with this situation in mind.

- Properly initialize replication slot state when recycling a previously-used slot (Michael Paquier)
  
  This failure to reset all of the fields of the slot could prevent VACUUM from removing dead tuples.

- Round shared-memory allocation request to a multiple of the actual huge page size when attempting to use huge pages on Linux (Tom Lane)
  
  This avoids possible failures during munmap() on systems with atypical default huge page sizes. Except in crash-recovery cases, there were no ill effects other than a log message.

- Use a more random value for the dynamic shared memory control segment's ID (Robert Haas, Tom Lane)
  
  Previously, the same value would be chosen every time, because it was derived from random() but srand() had not yet been called. While relatively harmless, this was not the intended behavior.

- On Windows, retry creation of the dynamic shared memory control segment after an access-denied error (Kyotaro Horiguchi, Amit Kapila)
  
  Windows sometimes returns ERROR_ACCESS_DENIED rather than ERROR_ALREADY_EXISTS when there is an existing segment. This led to postmaster startup failure due to believing that the former was an unrecoverable error.

- Fix PL/pgSQL to not misbehave with parameters and local variables of type int2vector or oidvector (Tom Lane)
• Don’t try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)

This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)

The reported problem involved leaking an error report during PQreset(), but there might be related cases.

• Make ecpg’s --help and --version options work consistently with our other executables (Haribabu Kommi)

• Fix pgbench’s calculation of average latency (Fabien Coelho)

The calculation was incorrect when there were \sleep commands in the script, or when the test duration was specified in number of transactions rather than total time.

• In pg_upgrade, check library loadability in name order (Tom Lane)

This is a workaround to deal with cross-extension dependencies from language transform modules to their base language and data type modules.

• In pg_dump, never dump range constructor functions (Tom Lane)

This oversight led to pg_upgrade failures with extensions containing range types, due to duplicate creation of the constructor functions.

• In pg_dump with -C, suppress TABLESPACE clause of CREATE DATABASE if --no-tablespaces is specified (Tom Lane)

• Make pg_receivexlog work correctly with --synchronous without slots (Gabriele Bartolini)

• Disallow specifying both --source-server and --source-target options to pg_rewind (Michael Banck)

• Make pg_rewind turn off synchronous_commit in its session on the source server (Michael Banck, Michael Paquier)

This allows pg_rewind to work even when the source server is using synchronous replication that is not working for some reason.

• In pg_xlogdump, retry opening new WAL segments when using --follow option (Magnus Hagander)

This allows for a possible delay in the server's creation of the next segment.

• Fix pg_xlogdump to cope with a WAL file that begins with a continuation record spanning more than one page (Pavan Deolasee)

• Fix contrib/pg_buffercache to work when shared_buffers exceeds 256GB (KaiGai Kohei)

• Fix contrib/intarray/bench/bench.pl to print the results of the EXPLAIN it does when given the -e option (Daniel Gustafsson)

• Support OpenSSL 1.1.0 (Heikki Linnakangas)

• Install TAP test infrastructure so that it’s available for extension testing (Craig Ringer)

When PostgreSQL has been configured with --enable-tap-tests, “make install” will now install the Perl support files for TAP testing where PGXS can find them. This allows non-core extensions to use $ (prove_check) without extra tests.
• In MSVC builds, include pg_recvlogical in a client-only installation (MauMau)

• Update Windows time zone mapping to recognize some time zone names added in recent Windows versions (Michael Paquier)

• Prevent failure of obsolete dynamic time zone abbreviations (Tom Lane)

If a dynamic time zone abbreviation does not match any entry in the referenced time zone, treat it as equivalent to the time zone name. This avoids unexpected failures when IANA removes abbreviations from their time zone database, as they did in tzdata release 2016f and seem likely to do again in the future. The consequences were not limited to not recognizing the individual abbreviation; any mismatch caused the pg_timezone_abbrevs view to fail altogether.

• Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.33. Release 9.5.4

Release date: 2016-08-11

This release contains a variety of fixes from 9.5.3. For information about new features in the 9.5 major release, see Section E.37.

E.33.1. Migration to Version 9.5.4

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.2, see Section E.35.

E.33.2. Changes

• Fix possible mis-evaluation of nested CASE-WHEN expressions (Heikki Linnakangas, Michael Paquier, Tom Lane)

A CASE expression appearing within the test value subexpression of another CASE could become confused about whether its own test value was null or not. Also, inlining of a SQL function implementing the equality operator used by a CASE expression could result in passing the wrong test value to functions called within a CASE expression in the SQL function's body. If the test values were of different data types, a crash might result; moreover such situations could be abused to allow disclosure of portions of server memory. (CVE-2016-5423)

• Fix client programs' handling of special characters in database and role names (Noah Misch, Nathan Bossart, Michael Paquier)

Numerous places in vacuumdb and other client programs could become confused by database and role names containing double quotes or backslashes. Tighten up quoting rules to make that safe. Also, ensure that when a conninfo string is used as a database name parameter to these programs, it is correctly treated as such throughout.
Fix handling of paired double quotes in psql's `\connect` and `\password` commands to match the documentation.

Introduce a new `-reuse-previous` option in psql's `\connect` command to allow explicit control of whether to re-use connection parameters from a previous connection. (Without this, the choice is based on whether the database name looks like a conninfo string, as before.) This allows secure handling of database names containing special characters in pg_dumpall scripts.

pg_dumpall now refuses to deal with database and role names containing carriage returns or new-lines, as it seems impractical to quote those characters safely on Windows. In future we may reject such names on the server side, but that step has not been taken yet.

These are considered security fixes because crafted object names containing special characters could have been used to execute commands with superuser privileges the next time a superuser executes pg_dumpall or other routine maintenance operations. (CVE-2016-5424)

• Fix corner-case misbehaviors for `IS NULL/IS NOT NULL` applied to nested composite values (Andrew Gierth, Tom Lane)

The SQL standard specifies that `IS NULL` should return `TRUE` for a row of all null values (thus `ROW(NULL,NULL) IS NULL` yields `TRUE`), but this is not meant to apply recursively (thus `ROW(NULL, ROW(NULL,NULL)) IS NULL` yields `FALSE`). The core executor got this right, but certain planner optimizations treated the test as recursive (thus producing `TRUE` in both cases), and contrib/postgres_fdw could produce remote queries that misbehaved similarly.

• Fix “unrecognized node type” error for `INSERT ... ON CONFLICT` within a recursive CTE (with item) (Peter Geoghegan)

• Fix `INSERT ... ON CONFLICT` to successfully match index expressions or index predicates that are simplified during the planner's expression preprocessing phase (Tom Lane)

• Correctly handle violations of exclusion constraints that apply to the target table of an `INSERT ... ON CONFLICT` command, but are not one of the selected arbiter indexes (Tom Lane)

    Such a case should raise a normal constraint-violation error, but it got into an infinite loop instead.

• Fix `INSERT ... ON CONFLICT` to not fail if the target table has a unique index on OID (Tom Lane)

• Make the `inet` and `cidr` data types properly reject IPv6 addresses with too many colon-separated fields (Tom Lane)

• Prevent crash in `close_ps()` (the point ## lseg operator) for NaN input coordinates (Tom Lane)

    Make it return NULL instead of crashing.

• Avoid possible crash in `pg_get_expr()` when inconsistent values are passed to it (Michael Paquier, Thomas Munro)

• Fix several one-byte buffer over-reads in `to_number()` (Peter Eisentraut)

    In several cases the `to_number()` function would read one more character than it should from the input string. There is a small chance of a crash, if the input happens to be adjacent to the end of memory.

• Do not run the planner on the query contained in `CREATE MATERIALIZED VIEW` or `CREATE TABLE AS` when `WITH NO DATA` is specified (Michael Paquier, Tom Lane)

    This avoids some unnecessary failure conditions, for example if a stable function invoked by the materialized view depends on a table that doesn’t exist yet.
• Avoid unsafe intermediate state during expensive paths through `heap_update()` (Masahiko Sawada, Andres Freund)

Previously, these cases locked the target tuple (by setting its XMAX) but did not WAL-log that action, thus risking data integrity problems if the page were spilled to disk and then a database crash occurred before the tuple update could be completed.

• Fix hint bit update during WAL replay of row locking operations (Andres Freund)

The only known consequence of this problem is that row locks held by a prepared, but uncommitted, transaction might fail to be enforced after a crash and restart.

• Avoid unnecessary “could not serialize access” errors when acquiring `FOR KEY SHARE` row locks in serializable mode (Álvaro Herrera)

• Make sure “expanded” datums returned by a plan node are read-only (Tom Lane)

This avoids failures in some cases where the result of a lower plan node is referenced in multiple places in upper nodes. So far as core PostgreSQL is concerned, only array values returned by PL/pgSQL functions are at risk; but extensions might use expanded datums for other things.

• Avoid crash in `postgres -C` when the specified variable has a null string value (Michael Paquier)

• Prevent unintended waits for the receiver in WAL sender processes (Kyotaro Horiguchi)

• Fix possible loss of large subtransactions in logical decoding (Petru-Florin Mihancea)

• Fix failure of logical decoding when a subtransaction contains no actual changes (Marko Tiikkaja, Andrew Gierth)

• Ensure that backends see up-to-date statistics for shared catalogs (Tom Lane)

The statistics collector failed to update the statistics file for shared catalogs after a request from a regular backend. This problem was partially masked because the autovacuum launcher regularly makes requests that did cause such updates; however, it became obvious with autovacuum disabled.

• Avoid redundant writes of the statistics files when multiple backends request updates close together (Tom Lane, Tomas Vondra)

• Avoid consuming a transaction ID during `VACUUM` (Alexander Korotkov)

Some cases in `VACUUM` unnecessarily caused an XID to be assigned to the current transaction. Normally this is negligible, but if one is up against the XID wraparound limit, consuming more XIDs during anti-wraparound vacuums is a very bad thing.

• Prevent possible failure when vacuuming multixact IDs in an installation that has been `pg_upgrade`d from pre-9.3 (Andrew Gierth, Álvaro Herrera)

The usual symptom of this bug is errors like “MultiXactId NNN has not been created yet -- apparent wraparound”.

• When a manual `ANALYZE` specifies a column list, don’t reset the table’s `changes_since_analyze` counter (Tom Lane)

If we’re only analyzing some columns, we should not prevent routine auto-analyze from happening for the other columns.

• Fix `ANALYZE`’s overestimation of `n_distinct` for a unique or nearly-unique column with many null entries (Tom Lane)

The nulls could get counted as though they were themselves distinct values, leading to serious planner misestimates in some types of queries.
• Prevent autovacuum from starting multiple workers for the same shared catalog (Álvaro Herrera)

Normally this isn't much of a problem because the vacuum doesn't take long anyway; but in the case of a severely bloated catalog, it could result in all but one worker uselessly waiting instead of doing useful work on other tables.

• Fix bug in b-tree mark/restore processing (Kevin Grittner)

This error could lead to incorrect join results or assertion failures in a merge join whose inner source node is a b-tree indexscan.

• Avoid duplicate buffer lock release when abandoning a b-tree index page deletion attempt (Tom Lane)

This mistake prevented VACUUM from completing in some cases involving corrupt b-tree indexes.

• Fix building of large (bigger than shared_buffers) hash indexes (Tom Lane)

The code path used for large indexes contained a bug causing incorrect hash values to be inserted into the index, so that subsequent index searches always failed, except for tuples inserted into the index after the initial build.

• Prevent infinite loop in GiST index build for geometric columns containing NaN component values (Tom Lane)

• Fix possible crash during a nearest-neighbor (ORDER BY distance) indexscan on a contrib/btree_gist index on an interval column (Peter Geoghegan)

• Fix “PANIC: failed to add BRIN tuple” error when attempting to update a BRIN index entry (Álvaro Herrera)

• Fix possible crash during background worker shutdown (Dmitry Ivanov)

• Fix PL/pgSQL’s handling of the INTO clause within IMPORT FOREIGN SCHEMA commands (Tom Lane)

• Fix contrib/btree_gin to handle the smallest possible bigint value correctly (Peter Eisentraut)

• Teach libpq to correctly decode server version from future servers (Peter Eisentraut)

It's planned to switch to two-part instead of three-part server version numbers for releases after 9.6. Make sure that PQserverVersion() returns the correct value for such cases.

• Fix ecpg’s code for unsigned long long array elements (Michael Meskes)

• In pg_dump with both -c and -C options, avoid emitting an unwanted CREATE SCHEMA public command (David Johnston, Tom Lane)

• Improve handling of SIGTERM/control-C in parallel pg_dump and pg_restore (Tom Lane)

Make sure that the worker processes will exit promptly, and also arrange to send query-cancel requests to the connected backends, in case they are doing something long-running such as a CREATE INDEX.

• Fix error reporting in parallel pg_dump and pg_restore (Tom Lane)

Previously, errors reported by pg_dump or pg_restore worker processes might never make it to the user's console, because the messages went through the master process, and there were various deadlock scenarios that would prevent the master process from passing on the messages. Instead, just print everything to stderr. In some cases this will result in duplicate messages (for instance, if all the workers report a server shutdown), but that seems better than no message.
• Ensure that parallel pg_dump or pg_restore on Windows will shut down properly after an error (Kyotaro Horiguchi)

Previously, it would report the error, but then just sit until manually stopped by the user.

• Make parallel pg_dump fail cleanly when run against a standby server (Magnus Hagander)

This usage is not supported unless --no-synchronized-snapshots is specified, but the error was not handled very well.

• Make pg_dump behave better when built without zlib support (Kyotaro Horiguchi)

It didn't work right for parallel dumps, and emitted some rather pointless warnings in other cases.

• Make pg_basebackup accept -Z 0 as specifying no compression (Fujii Masao)

• Fix makefiles' rule for building AIX shared libraries to be safe for parallel make (Noah Misch)

• Fix TAP tests and MSVC scripts to work when build directory's path name contains spaces (Michael Paquier, Kyotaro Horiguchi)

• Be more predictable about reporting “statement timeout” versus “lock timeout” (Tom Lane)

On heavily loaded machines, the regression tests sometimes failed due to reporting “lock timeout” even though the statement timeout should have occurred first.

• Make regression tests safe for Danish and Welsh locales (Jeff Janes, Tom Lane)

Change some test data that triggered the unusual sorting rules of these locales.

• Update our copy of the timezone code to match IANA's tzcode release 2016c (Tom Lane)

This is needed to cope with anticipated future changes in the time zone data files. It also fixes some corner-case bugs in coping with unusual time zones.

• Update time zone data files to tzdata release 2016f for DST law changes in Kemerovo and Novosibirsk, plus historical corrections for Azerbaijan, Belarus, and Morocco.

E.34. Release 9.5.3

Release date: 2016-05-12

This release contains a variety of fixes from 9.5.2. For information about new features in the 9.5 major release, see Section E.37.

E.34.1. Migration to Version 9.5.3

A dump/restore is not required for those running 9.5.X.

However, if you are upgrading from a version earlier than 9.5.2, see Section E.35.

E.34.2. Changes

• Clear the OpenSSL error queue before OpenSSL calls, rather than assuming it's clear already; and make sure we leave it clear afterwards (Peter Geoghegan, Dave Vitek, Peter Eisentraut)

This change prevents problems when there are multiple connections using OpenSSL within a single process and not all the code involved follows the same rules for when to clear the error queue. Failures have been reported specifically when a client application uses SSL connections in libpq...
concurrently with SSL connections using the PHP, Python, or Ruby wrappers for OpenSSL. It's possible for similar problems to arise within the server as well, if an extension module establishes an outgoing SSL connection.

- Fix “failed to build any n-way joins” planner error with a full join enclosed in the right-hand side of a left join (Tom Lane)

- Fix incorrect handling of equivalence-class tests in multilevel nestloop plans (Tom Lane)

  Given a three-or-more-way equivalence class of variables, such as \(X.X = Y.Y = Z.Z\), it was possible for the planner to omit some of the tests needed to enforce that all the variables are actually equal, leading to join rows being output that didn't satisfy the \(\text{WHERE}\) clauses. For various reasons, erroneous plans were seldom selected in practice, so that this bug has gone undetected for a long time.

- Fix corner-case parser failures occurring when \texttt{operator\_precedence\_warning} is turned on (Tom Lane)

  An example is that \texttt{SELECT (ARRAY[])}::text[] gave an error, though it worked without the parentheses.

- Fix query-lifespan memory leak in GIN index scans (Julien Rouhaud)

- Fix query-lifespan memory leak and potential index corruption hazard in GIN index insertion (Tom Lane)

  The memory leak would typically not amount to much in simple queries, but it could be very substantial during a large GIN index build with high \texttt{maintenance\_work\_mem}.

- Fix possible misbehavior of \texttt{TH}, \texttt{th}, and \(Y, YYY\) format codes in \texttt{to\_timestamp()} (Tom Lane)

  These could advance off the end of the input string, causing subsequent format codes to read garbage.

- Fix dumping of rules and views in which the \texttt{array} argument of a \texttt{value operator ANY (array)} construct is a sub-\texttt{SELECT} (Tom Lane)

- Disallow newlines in \texttt{ALTER SYSTEM} parameter values (Tom Lane)

  The configuration-file parser doesn't support embedded newlines in string literals, so we mustn't allow them in values to be inserted by \texttt{ALTER SYSTEM}.

- Fix \texttt{ALTER TABLE ... REPLICA IDENTITY USING INDEX} to work properly if an index on OID is selected (David Rowley)

- Avoid possible misbehavior after failing to remove a tablespace symlink (Tom Lane)

- Fix crash in logical decoding on alignment-picky platforms (Tom Lane, Andres Freund)

  The failure occurred only with a transaction large enough to spill to disk and a primary-key change within that transaction.

- Avoid repeated requests for feedback from receiver while shutting down walsender (Nick Cleaton)

- Make \texttt{pg_regress} use a startup timeout from the \texttt{PGCTLTIMEOUT} environment variable, if that's set (Tom Lane)

  This is for consistency with a behavior recently added to \texttt{pg\_ctl}; it eases automated testing on slow machines.

- Fix \texttt{pg\_upgrade} to correctly restore extension membership for operator families containing only one operator class (Tom Lane)
In such a case, the operator family was restored into the new database, but it was no longer marked as part of the extension. This had no immediate ill effects, but would cause later pg_dump runs to emit output that would cause (harmless) errors on restore.

- Fix pg_upgrade to not fail when new-cluster TOAST rules differ from old (Tom Lane)

pg_upgrade had special-case code to handle the situation where the new PostgreSQL version thinks that a table should have a TOAST table while the old version did not. That code was broken, so remove it, and instead do nothing in such cases; there seems no reason to believe that we can't get along fine without a TOAST table if that was okay according to the old version's rules.

- Fix atomic operations for PPC when using IBM's XLC compiler (Noah Misch)

- Reduce the number of SysV semaphores used by a build configured with --disable-spin-locks (Tom Lane)

- Rename internal function strtoi() to strtoint() to avoid conflict with a NetBSD library function (Thomas Munro)

- Fix reporting of errors from bind() and listen() system calls on Windows (Tom Lane)

- Reduce verbosity of compiler output when building with Microsoft Visual Studio (Christian Ullrich)

- Support building with Visual Studio 2015 (Michael Paquier, Petr Jelínek)

  Note that builds made with VS2015 will not run on Windows versions before Windows Vista.

- Fix putenv() to work properly with Visual Studio 2013 (Michael Paquier)

- Avoid possibly-unsafe use of Windows' FormatMessage() function (Christian Ullrich)

  Use the FORMAT_MESSAGE_IGNORE_INSERTS flag where appropriate. No live bug is known to exist here, but it seems like a good idea to be careful.

- Update time zone data files to tzdata release 2016d for DST law changes in Russia and Venezuela. There are new zone names Europe/Kirov and Asia/Tomsk to reflect the fact that these regions now have different time zone histories from adjacent regions.

### E.35. Release 9.5.2

**Release date:** 2016-03-31

This release contains a variety of fixes from 9.5.1. For information about new features in the 9.5 major release, see Section E.37.

#### E.35.1. Migration to Version 9.5.2

A dump/restore is not required for those running 9.5.X.

However, you may need to REINDEX some indexes after applying the update, as per the first changelog entry below.

#### E.35.2. Changes

- Disable abbreviated keys for string sorting in non-C locales (Robert Haas)

  PostgreSQL 9.5 introduced logic for speeding up comparisons of string data types by using the standard C library function strxfrm() as a substitute for strcoll(). It now emerges that most versions of glibc (Linux's implementation of the C library) have buggy implementations of
strxfrm() that, in some locales, can produce string comparison results that do not match strcoll(). Until this problem can be better characterized, disable the optimization in all non-C locales. (C locale is safe since it uses neither strcoll() nor strxfrm().)

Unfortunately, this problem affects not only sorting but also entry ordering in B-tree indexes, which means that B-tree indexes on text, varchar, or char columns may now be corrupt if they sort according to an affected locale and were built or modified under PostgreSQL 9.5.0 or 9.5.1. Users should REINDEX indexes that might be affected.

It is not possible at this time to give an exhaustive list of known-affected locales. C locale is known safe, and there is no evidence of trouble in English-based locales such as en_US, but some other popular locales such as de_DE are affected in most glibc versions.

- Maintain row-security status properly in cached plans (Stephen Frost)

> In a session that performs queries as more than one role, the plan cache might incorrectly re-use a plan that was generated for another role ID, thus possibly applying the wrong set of policies when row-level security (RLS) is in use. (CVE-2016-2193)

- Add must-be-superuser checks to some new contrib/pageinspect functions (Andreas Seltenreich)

> Most functions in the pageinspect extension that inspect bytea values disallow calls by non-superusers, but brin_page_type() and brin_metapage_info() failed to do so. Passing contrived bytea values to them might crash the server or disclose a few bytes of server memory. Add the missing permissions checks to prevent misuse. (CVE-2016-3065)

- Fix incorrect handling of indexed \texttt{ROW()} comparisons (Simon Riggs)

Flaws in a minor optimization introduced in 9.5 caused incorrect results if the \texttt{ROW()} comparison matches the index ordering partially but not exactly (for example, differing column order, or the index contains both ASC and DESC columns). Pending a better solution, the optimization has been removed.

- Fix incorrect handling of NULL index entries in indexed \texttt{ROW()} comparisons (Tom Lane)

An index search using a row comparison such as \texttt{ROW(a, b) > ROW('x', 'y')} would stop upon reaching a NULL entry in the b column, ignoring the fact that there might be non-NULL b values associated with later values of a.

- Avoid unlikely data-loss scenarios due to renaming files without adequate \texttt{fsync()} calls before and after (Michael Paquier, Tomas Vondra, Andres Freund)

- Fix incorrect behavior when rechecking a just-modified row in a query that does SELECT FOR UPDATE/SHARE and contains some relations that need not be locked (Tom Lane)

Rows from non-locked relations were incorrectly treated as containing all NULLs during the recheck, which could result in incorrectly deciding that the updated row no longer passes the \texttt{WHERE} condition, or in incorrectly outputting NULLs.

- Fix bug in \texttt{json_to_record()} when a field of its input object contains a sub-object with a field name matching one of the requested output column names (Tom Lane)

- Fix nonsense result from two-argument form of \texttt{jsonb_object()} when called with empty arrays (Michael Paquier, Andrew Dunstan)

- Fix misbehavior in \texttt{jsonb_set()} when converting a path array element into an integer for use as an array subscript (Michael Paquier)

- Fix misformatting of negative time zone offsets by \texttt{to_char()}’s \texttt{OF} format code (Thomas Munro, Tom Lane)

2523
• Fix possible incorrect logging of waits done by INSERT ... ON CONFLICT (Peter Geoghegan)
  Log messages would sometimes claim that the wait was due to an exclusion constraint although no such constraint was responsible.

• Ignore recovery_min_apply_delay parameter until recovery has reached a consistent state (Michael Paquier)
  Previously, standby servers would delay application of WAL records in response to recovery_min_apply_delay even while replaying the initial portion of WAL needed to make their database state valid. Since the standby is useless until it's reached a consistent database state, this was deemed unhelpful.

• Correctly handle cases where pg_subtrans is close to XID wraparound during server startup (Jeff Janes)

• Fix assorted bugs in logical decoding (Andres Freund)
  Trouble cases included tuples larger than one page when replica identity is FULL, UPDATEs that change a primary key within a transaction large enough to be spooled to disk, incorrect reports of “subxact logged without previous toplevel record”, and incorrect reporting of a transaction’s commit time.

• Fix planner error with nested security barrier views when the outer view has a WHERE clause containing a correlated subquery (Dean Rasheed)

• Fix memory leak in GIN index searches (Tom Lane)

• Fix corner-case crash due to trying to free localeconv() output strings more than once (Tom Lane)

• Fix parsing of affix files for ispell dictionaries (Tom Lane)
  The code could go wrong if the affix file contained any characters whose byte length changes during case-folding, for example I in Turkish UTF8 locales.

• Avoid use of sscanf() to parse ispell dictionary files (Artur Zakirov)
  This dodges a portability problem on FreeBSD-derived platforms (including macOS).

• Fix atomic-operations code used on PPC with IBM's xlc compiler (Noah Misch)
  This error led to rare failures of concurrent operations on that platform.

• Avoid a crash on old Windows versions (before 7SP1/2008R2SP1) with an AVX2-capable CPU and a Postgres build done with Visual Studio 2013 (Christian Ullrich)
  This is a workaround for a bug in Visual Studio 2013’s runtime library, which Microsoft have stated they will not fix in that version.

• Fix psql's tab completion logic to handle multibyte characters properly (Kyotaro Horiguchi, Robert Haas)

• Fix psql's tab completion for SECURITY LABEL (Tom Lane)
  Pressing TAB after SECURITY LABEL might cause a crash or offering of inappropriate keywords.

• Make pg_ctl accept a wait timeout from the PGCTLTIMEOUT environment variable, if none is specified on the command line (Noah Misch)
  This eases testing of slower buildfarm members by allowing them to globally specify a longer-than-normal timeout for postmaster startup and shutdown.
• Fix incorrect test for Windows service status in pg_ctl (Manuel Mathar)

The previous set of minor releases attempted to fix pg_ctl to properly determine whether to send log messages to Window's Event Log, but got the test backwards.

• Fix pgbench to correctly handle the combination of -c and -M prepared options (Tom Lane)

• In pg_upgrade, skip creating a deletion script when the new data directory is inside the old data directory (Bruce Momjian)

  Blind application of the script in such cases would result in loss of the new data directory.

• In PL/Perl, properly translate empty Postgres arrays into empty Perl arrays (Alex Hunsaker)

• Make PL/Python cope with function names that aren't valid Python identifiers (Jim Nasby)

• Fix multiple mistakes in the statistics returned by contrib/pgstattuple's pgstatindex() function (Tom Lane)

• Remove dependency on pseed in MSVC builds, since it's no longer provided by core Perl (Michael Paquier, Andrew Dunstan)

• Update time zone data files to tzdata release 2016c for DST law changes in Azerbaijan, Chile, Haiti, Palestine, and Russia (Altai, Astrakhan, Kirov, Sakhalin, Ulyanovsk regions), plus historical corrections for Lithuania, Moldova, and Russia (Kaliningrad, Samara, Volgograd).

E.36. Release 9.5.1

Release date: 2016-02-11

This release contains a variety of fixes from 9.5.0. For information about new features in the 9.5 major release, see Section E.37.

E.36.1. Migration to Version 9.5.1

A dump/restore is not required for those running 9.5.X.

E.36.2. Changes

• Fix infinite loops and buffer-overrun problems in regular expressions (Tom Lane)

  Very large character ranges in bracket expressions could cause infinite loops in some cases, and memory overwrites in other cases. (CVE-2016-0773)

• Fix an oversight that caused hash joins to miss joining to some tuples of the inner relation in rare cases (Tomas Vondra, Tom Lane)

• Avoid pushdown of HAVING clauses when grouping sets are used (Andrew Gierth)

• Fix deparsing of ON CONFLICT arbiter WHERE clauses (Peter Geoghegan)

• Make %h and %r escapes in log_line_prefix work for messages emitted due to log_connections (Tom Lane)

  Previously, %h/%r started to work just after a new session had emitted the “connection received” log message; now they work for that message too.

• Avoid leaking a token handle during SSPI authentication (Christian Ullrich)
• Fix psql’s `\d` command to interpret its pattern argument the same way as other `\d` commands with potentially schema-qualified patterns do (Reece Hart)

• In pg_ctl on Windows, check service status to decide where to send output, rather than checking if standard output is a terminal (Michael Paquier)

• Fix assorted corner-case bugs in pg_dump’s processing of extension member objects (Tom Lane)

• Fix improper quoting of domain constraint names in pg_dump (Elvis Pranskevichus)

• Make pg_dump mark a view’s triggers as needing to be processed after its rule, to prevent possible failure during parallel pg_restore (Tom Lane)

• Install guards in pgbench against corner-case overflow conditions during evaluation of script-specified division or modulo operators (Fabien Coelho, Michael Paquier)

• Suppress useless warning message when pg_receivexlog connects to a pre-9.4 server (Marco Nenciarini)

• Avoid dump/reload problems when using both plpython2 and plpython3 (Tom Lane)

  In principle, both versions of PL/Python can be used in the same database, though not in the same session (because the two versions of libpython cannot safely be used concurrently). However, pg_restore and pg_upgrade both do things that can fall foul of the same-session restriction. Work around that by changing the timing of the check.

• Fix PL/Python regression tests to pass with Python 3.5 (Peter Eisentraut)

• Prevent certain PL/Java parameters from being set by non-superusers (Noah Misch)

  This change mitigates a PL/Java security bug (CVE-2016-0766), which was fixed in PL/Java by marking these parameters as superuser-only. To fix the security hazard for sites that update PostgreSQL more frequently than PL/Java, make the core code aware of them also.

• Fix ecpg-supplied header files to not contain comments continued from a preprocessor directive line onto the next line (Michael Meskes)

  Such a comment is rejected by ecpg. It’s not yet clear whether ecpg itself should be changed.

• Fix `hstore_to_json_loose()`’s test for whether an hstore value can be converted to a JSON number (Tom Lane)

  Previously this function could be fooled by non-alphanumeric trailing characters, leading to emitting syntactically-invalid JSON.

• In contrib/postgres_fdw, fix bugs triggered by use of `tableoid` in data-modifying commands (Etsuro Fujita, Robert Haas)

• Fix ill-advised restriction of NAMEDATALEN to be less than 256 (Robert Haas, Tom Lane)

• Improve reproducibility of build output by ensuring filenames are given to the linker in a fixed order (Christoph Berg)

  This avoids possible bitwise differences in the produced executable files from one build to the next.

• Ensure that `dynloader.h` is included in the installed header files in MSVC builds (Bruce Momjian, Michael Paquier)

• Update time zone data files to tzdata release 2016a for DST law changes in Cayman Islands, Metlakatla, and Trans-Baikal Territory (Zabaykalsky Krai), plus historical corrections for Pakistan.
E.37. Release 9.5

Release date: 2016-01-07

E.37.1. Overview

Major enhancements in PostgreSQL 9.5 include:

- Allow **INSERT**s that would generate constraint conflicts to be turned into **UPDATE**s or ignored
- Add **GROUP BY** analysis features **GROUPING SETS**, **CUBE** and **ROLLUP**
- Add row-level security control
- Create mechanisms for tracking the progress of replication, including methods for identifying the origin of individual changes during logical replication
- Add Block Range Indexes (BRIN)
- Substantial performance improvements for sorting
- Substantial performance improvements for multi-CPU machines

The above items are explained in more detail in the sections below.

E.37.2. Migration to Version 9.5

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.5 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

- Adjust operator precedence to match the SQL standard (Tom Lane)

  The precedence of <=, >= and <> has been reduced to match that of <, > and =. The precedence of **IS** tests (e.g., x **IS** NULL) has been reduced to be just below these six comparison operators. Also, multi-keyword operators beginning with **NOT** now have the precedence of their base operator (for example, **NOT BETWEEN** now has the same precedence as **BETWEEN** whereas before they had inconsistent precedence, behaving like **NOT** with respect to their left operand but like their base operator with respect to their right operand. The new configuration parameter **operator_precedence_warning** can be enabled to warn about queries in which these precedence changes result in different parsing choices.

- Change pg_ctl’s default shutdown mode from **smart** to **fast** (Bruce Momjian)

  This means the default behavior will be to forcibly cancel existing database sessions, not simply wait for them to exit.

- Use assignment cast behavior for data type conversions in PL/pgSQL assignments, rather than converting to and from text (Tom Lane)

  This change causes conversions of Booleans to strings to produce **true** or **false**, not t or f. Other type conversions may succeed in more cases than before; for example, assigning a numeric value 3.9 to an integer variable will now assign 4 rather than failing. If no assignment-grade cast is defined for the particular source and destination types, PL/pgSQL will fall back to its old I/O conversion behavior.

- Allow characters in server command-line options to be escaped with a backslash (Andres Freund)
Formerly, spaces in the options string always separated options, so there was no way to include a space in an option value. Including a backslash in an option value now requires writing \\

- Change the default value of the GSSAPI include_realm parameter to 1, so that by default the realm is not removed from a GSS or SSPI principal name (Stephen Frost)

- Replace configuration parameter checkpoint_segments with min_wal_size and max_wal_size (Heikki Linnakangas)

If you previously adjusted checkpoint_segments, the following formula will give you an approximately equivalent setting:

\[
max\_wal\_size = (3 \times \text{checkpoint\_segments}) \times 16\text{MB}
\]

Note that the default setting for max_wal_size is much higher than the default checkpoint_segments used to be, so adjusting it might no longer be necessary.

- Control the Linux OOM killer via new environment variables PG_OOM_ADJUST_FILE and PG_OOM_ADJUST_VALUE, instead of compile-time options LINUX_OOM_SCORE_ADJ and LINUX_OOM_ADJ (Gurjeet Singh)

- Decommission server configuration parameter ssl_renegotiation_limit, which was deprecated in earlier releases (Andres Freund)

While SSL renegotiation is a good idea in theory, it has caused enough bugs to be considered a net negative in practice, and it is due to be removed from future versions of the relevant standards. We have therefore removed support for it from PostgreSQL. The ssl_renegotiation_limit parameter still exists, but cannot be set to anything but zero (disabled). It's not documented anymore, either.

- Remove server configuration parameter autocommit, which was already deprecated and non-operational (Tom Lane)

- Remove the pg_authid catalog's rolcatupdate field, as it had no usefulness (Adam Brightwell)

- The pg_stat_replication system view's sent field is now NULL, not zero, when it has no valid value (Magnus Hagander)

- Allow json and jsonb array extraction operators to accept negative subscripts, which count from the end of JSON arrays (Peter Geoghegan, Andrew Dunstan)

Previously, these operators returned NULL for negative subscripts.

**E.37.3. Changes**

Below you will find a detailed account of the changes between PostgreSQL 9.5 and the previous major release.

**E.37.3.1. Server**

**E.37.3.1.1. Indexes**

- Add Block Range Indexes (BRIN) (Álvaro Herrera)

BRIN indexes store only summary data (such as minimum and maximum values) for ranges of heap blocks. They are therefore very compact and cheap to update; but if the data is naturally clustered, they can still provide substantial speedup of searches.
• Allow queries to perform accurate distance filtering of bounding-box-indexed objects (polygons, circles) using GiST indexes (Alexander Korotkov, Heikki Linnakangas)

Previously, to exploit such an index a subquery had to be used to select a large number of rows ordered by bounding-box distance, and the result then had to be filtered further with a more accurate distance calculation.

• Allow GiST indexes to perform index-only scans (Anastasia Lubennikova, Heikki Linnakangas, Andreas Karlsson)

• Add configuration parameter gin_pending_list_limit to control the size of GIN pending lists (Fujii Masao)

This value can also be set on a per-index basis as an index storage parameter. Previously the pending-list size was controlled by work_mem, which was awkward because appropriate values for work_mem are often much too large for this purpose.

• Issue a warning during the creation of hash indexes because they are not crash-safe (Bruce Momjian)

E.37.3.1.2. General Performance

• Improve the speed of sorting of varchar, text, and numeric fields via “abbreviated” keys (Peter Geoghegan, Andrew Gierth, Robert Haas)

• Extend the infrastructure that allows sorting to be performed by inlined, non-SQL-callable comparison functions to cover CREATE INDEX, REINDEX, and CLUSTER (Peter Geoghegan)

• Improve performance of hash joins (Tomas Vondra, Robert Haas)

• Improve concurrency of shared buffer replacement (Robert Haas, Amit Kapila, Andres Freund)

• Reduce the number of page locks and pins during index scans (Kevin Grittner)

  The primary benefit of this is to allow index vacuums to be blocked less often.

• Make per-backend tracking of buffer pins more memory-efficient (Andres Freund)

• Improve lock scalability (Andres Freund)

  This particularly addresses scalability problems when running on systems with multiple CPU sockets.

• Allow the optimizer to remove unnecessary references to left-joined subqueries (David Rowley)

• Allow pushdown of query restrictions into subqueries with window functions, where appropriate (David Rowley)

• Allow a non-leakproof function to be pushed down into a security barrier view if the function does not receive any view output columns (Dean Rasheed)

• Teach the planner to use statistics obtained from an expression index on a boolean-returning function, when a matching function call appears in WHERE (Tom Lane)

• Make ANALYZE compute basic statistics (null fraction and average column width) even for columns whose data type lacks an equality function (Oleksandr Shulgin)

• Speed up CRC (cyclic redundancy check) computations and switch to CRC-32C (Abhijit Menon-Sen, Heikki Linnakangas)

• Improve bitmap index scan performance (Teodor Sigaev, Tom Lane)

• Speed up CREATE INDEX by avoiding unnecessary memory copies (Robert Haas)
• Increase the number of buffer mapping partitions (Amit Kapila, Andres Freund, Robert Haas)

This improves performance for highly concurrent workloads.

E.37.3.1.3. Monitoring

• Add per-table autovacuum logging control via new log_autovacuum_min_duration storage parameter (Michael Paquier)

• Add new configuration parameter cluster_name (Thomas Munro)

This string, typically set in postgresql.conf, allows clients to identify the cluster. This name also appears in the process title of all server processes, allowing for easier identification of processes belonging to the same cluster.

• Prevent non-superusers from changing log_disconnections on connection startup (Fujii Masao)

E.37.3.1.4. SSL

• Check “Subject Alternative Names” in SSL server certificates, if present (Alexey Klyukin)

When they are present, this replaces checks against the certificate’s “Common Name”.

• Add system view pg_stat_ssl to report SSL connection information (Magnus Hagander)

• Add libpq functions to return SSL information in an implementation-independent way (Heikki Linnakangas)

While PQgetssl() can still be used to call OpenSSL functions, it is now considered deprecated because future versions of libpq might support other SSL implementations. When possible, use the new functions PQsslAttribute(), PQsslAttributeNames(), and PQsslInUse() to obtain SSL information in an SSL-implementation-independent way.

• Make libpq honor any OpenSSL thread callbacks (Jan Urbanski)

Previously they were overwritten.

E.37.3.1.5. Server Settings

• Replace configuration parameter checkpoint_segments with min_wal_size and max_wal_size (Heikki Linnakangas)

This change allows the allocation of a large number of WAL files without keeping them after they are no longer needed. Therefore the default for max_wal_size has been set to 1GB, much larger than the old default for checkpoint_segments. Also note that standby servers perform restartpoints to try to limit their WAL space consumption to max_wal_size; previously they did not pay any attention to checkpoint_segments.

• Control the Linux OOM killer via new environment variables PG_OOM_ADJUST_FILE and PG_OOM_ADJUST_VALUE (Gurjeet Singh)

The previous OOM control infrastructure involved compile-time options LINUX_OOM_S-CORE_ADJ and LINUX_OOM_ADJ, which are no longer supported. The new behavior is available in all builds.

• Allow recording of transaction commit time stamps when configuration parameter track_commit_timestamp is enabled (Álvaro Herrera, Petr Jelínek)

Time stamp information can be accessed using functions pg_xact_commit_timestamp() and pg_last_committed_xact().
• Allow `local_preload_libraries` to be set by `ALTER ROLE SET` (Peter Eisentraut, Kyotaro Horiguchi)

• Allow autovacuum workers to respond to configuration parameter changes during a run (Michael Paquier)

• Make configuration parameter `debug_assertions` read-only (Andres Freund)

  This means that assertions can no longer be turned off if they were enabled at compile time, allowing for more efficient code optimization. This change also removes the `postgres -A` option.

• Allow setting `effective_io_concurrency` on systems where it has no effect (Peter Eisentraut)

• Add system view `pg_file_settings` to show the contents of the server's configuration files (Sawada Masahiko)

• Add `pending_restart` to the system view `pg_settings` to indicate a change has been made but will not take effect until a database restart (Peter Eisentraut)

• Allow `ALTER SYSTEM` values to be reset with `ALTER SYSTEM RESET` (Vik Fearing)

  This command removes the specified setting from `postgresql.auto.conf`.

E.37.3.2. Replication and Recovery

• Create mechanisms for tracking the progress of replication, including methods for identifying the origin of individual changes during logical replication (Andres Freund)

  This is helpful when implementing replication solutions.

• Rework truncation of the multixact commit log to be properly WAL-logged (Andres Freund)

  This makes things substantially simpler and more robust.

• Add `recovery.conf` parameter `recovery_target_action` to control post-recovery activity (Petr Jelínek)

  This replaces the old parameter `pause_at_recovery_target`.

• Add new archive_mode value `always` to allow standbys to always archive received WAL files (Fujii Masao)

• Add configuration parameter `wal_retrieve_retry_interval` to control WAL read retry after failure (Alexey Vasiliev, Michael Paquier)

  This is particularly helpful for warm standbys.

• Allow compression of full-page images stored in WAL (Rahila Syed, Michael Paquier)

  This feature reduces WAL volume, at the cost of more CPU time spent on WAL logging and WAL replay. It is controlled by a new configuration parameter `wal_compression`, which currently is off by default.

• Archive WAL files with suffix `.partial` during standby promotion (Heikki Linnakangas)

• Add configuration parameter `log_replication_commands` to log replication commands (Fujii Masao)

  By default, replication commands, e.g. `IDENTIFY_SYSTEM`, are not logged, even when `log_statement` is set to `all`.

• Report the processes holding replication slots in `pg_replication_slots` (Craig Ringer)
The new output column is `active_pid`.

- Allow `recovery.conf`'s `primary_conninfo` setting to use connection URIs, e.g. `postgres://` (Alexander Shulgin)

E.37.3.3. Queries

- Allow `INSERT`s that would generate constraint conflicts to be turned into `UPDATE`s or ignored (Peter Geoghegan, Heikki Linnakangas, Andres Freund)

  The syntax is `INSERT ... ON CONFLICT DO NOTHING/UPDATE`. This is the Postgres implementation of the popular `UPSERT` command.

- Add `GROUP BY` analysis features `GROUPING SETS`, `CUBE` and `ROLLUP` (Andrew Gierth, Atri Sharma)

- Allow setting multiple target columns in an `UPDATE` from the result of a single sub-SELECT (Tom Lane)

  This is accomplished using the syntax `UPDATE tab SET (col1, col2, ...) = (SELECT ...)`.

- Add `SELECT` option `SKIP LOCKED` to skip locked rows (Thomas Munro)

  This does not throw an error for locked rows like `NOWAIT` does.

- Add `SELECT` option `TABLESAMPLE` to return a subset of a table (Petr Jelínek)

  This feature supports the SQL-standard table sampling methods. In addition, there are provisions for user-defined table sampling methods.

- Suggest possible matches for mistyped column names (Peter Geoghegan, Robert Haas)

E.37.3.4. Utility Commands

- Add more details about sort ordering in `EXPLAIN` output (Marius Timmer, Lukas Kreft, Arne Scheffer)

  Details include `COLLATE`, `DESC`, `USING`, and `NULLS FIRST/LAST`.

- Make `VACUUM` log the number of pages skipped due to pins (Jim Nasby)

- Make `TRUNCATE` properly update the `pg_stat*` tuple counters (Alexander Shulgin)

E.37.3.4.1. REINDEX

- Allow `REINDEX` to reindex an entire schema using the `SCHEMA` option (Sawada Masahiko)

- Add `VERBOSE` option to `REINDEX` (Sawada Masahiko)

- Prevent `REINDEX DATABASE` and `SCHEMA` from outputting object names, unless `VERBOSE` is used (Simon Riggs)

- Remove obsolete `FORCE` option from `REINDEX` (Fujii Masao)

E.37.3.5. Object Manipulation

- Add row-level security control (Craig Ringer, KaiGai Kohei, Adam Brightwell, Dean Rasheed, Stephen Frost)
This feature allows row-by-row control over which users can add, modify, or even see rows in a table. This is controlled by new commands CREATE/ALTER/DROP POLICY and ALTER TABLE ... ENABLE/DISABLE ROW SECURITY.

- Allow changing of the WAL logging status of a table after creation with ALTER TABLE ... SET LOGGED / UNLOGGED (Fabrízio de Royes Mello)
- Add IF NOT EXISTS clause to CREATE TABLE AS, CREATE INDEX, CREATE SEQUENCE, and CREATE MATERIALIZED VIEW (Fabrízio de Royes Mello)
- Add support for IF EXISTS to ALTER TABLE ... RENAME CONSTRAINT (Bruce Momjian)
- Allow some DDL commands to accept CURRENT_USER or SESSION_USER, meaning the current user or session user, in place of a specific user name (Kyotaro Horiguchi, Álvaro Herrera)

This feature is now supported in ALTER USER, ALTER GROUP, ALTER ROLE, GRANT, and ALTER object OWNER TO commands.

- Support comments on domain constraints (Álvaro Herrera)
- Reduce lock levels of some create and alter trigger and foreign key commands (Simon Riggs, Andreas Karlsson)
- Allow LOCK TABLE ... ROW EXCLUSIVE MODE for those with INSERT privileges on the target table (Stephen Frost)

Previously this command required UPDATE, DELETE, or TRUNCATE privileges.

- Apply table and domain CHECK constraints in order by name (Tom Lane)

The previous ordering was indeterminate.

- Allow CREATE/ALTER DATABASE to manipulate datistemplate and datallowconn (Vik Fearing)

This allows these per-database settings to be changed without manually modifying the pg_data-base system catalog.

### E.37.3.5.1. Foreign Tables

- Add support for IMPORT FOREIGN SCHEMA (Ronan Dunklau, Michael Paquier, Tom Lane)

This command allows automatic creation of local foreign tables that match the structure of existing tables on a remote server.

- Allow CHECK constraints to be placed on foreign tables (Shigeru Hanada, Etsuro Fujita)

Such constraints are assumed to be enforced on the remote server, and are not enforced locally. However, they are assumed to hold for purposes of query optimization, such as constraint exclusion.

- Allow foreign tables to participate in inheritance (Shigeru Hanada, Etsuro Fujita)

To let this work naturally, foreign tables are now allowed to have check constraints marked as not valid, and to set storage and OID characteristics, even though these operations are effectively no-ops for a foreign table.

- Allow foreign data wrappers and custom scans to implement join pushdown (KaiGai Kohei)

### E.37.3.5.2. Event Triggers

- Whenever a ddl_command_end event trigger is installed, capture details of DDL activity for it to inspect (Álvaro Herrera)
This information is available through a set-returning function `pg_event_trigger_ddl_commands()`, or by inspection of C data structures if that function doesn't provide enough detail.

- Allow event triggers on table rewrites caused by `ALTER TABLE` (Dimitri Fontaine)
- Add event trigger support for database-level `COMMENT`, `SECURITY LABEL`, and `GRANT/REVOKE` (Álvaro Herrera)
- Add columns to the output of `pg_event_trigger_dropped_objects` (Álvaro Herrera)

This allows simpler processing of delete operations.

### E.37.3.6. Data Types

- Allow the `xml` data type to accept empty or all-whitespace content values (Peter Eisentraut)
  
  This is required by the SQL/XML specification.
- Allow `macaddr` input using the format `xxxx-xxxx-xxxx` (Herwin Weststrate)
- Disallow non-SQL-standard syntax for `interval` with both precision and field specifications (Bruce Momjian)

  Per the standard, such type specifications should be written as, for example, `INTERVAL MINUTE TO SECOND(2)`.

  PostgreSQL formerly allowed this to be written as `INTERVAL(2) MINUTE TO SECOND`, but it must now be written in the standard way.
- Add selectivity estimators for `inet/cidr` operators and improve estimators for text search functions (Emre Hasegeli, Tom Lane)
- Add data types `regrole` and `regnamespace` to simplify entering and pretty-printing the OID of a role or namespace (Kyotaro Horiguchi)

### E.37.3.6.1. JSON

- Add `jsonb` functions `jsonb_set()` and `jsonb_pretty()` (Dmitry Dolgov, Andrew Dunstan, Petr Jelinek)
- Add `jsonb` generator functions `to_jsonb()`, `jsonb_object()`, `jsonb_object()`, `jsonb_build_object()`, `jsonb_build_array()`, `jsonb_agg()`, and `jsonb_object_agg()` (Andrew Dunstan)

  Equivalent functions already existed for type `json`.
- Reduce casting requirements `to/from json and jsonb` (Tom Lane)
- Allow `text`, `text array`, and `integer` values to be subtracted from `jsonb` documents (Dmitry Dolgov, Andrew Dunstan)
- Add `jsonb ||` operator (Dmitry Dolgov, Andrew Dunstan)
- Add `json_strip_nulls()` and `jsonb_strip_nulls()` functions to remove JSON null values from documents (Andrew Dunstan)

### E.37.3.7. Functions

- Add `generate_series()` for numeric values (Plato Malugin)
- Allow `array_agg()` and `ARRAY()` to take arrays as inputs (Ali Akbar, Tom Lane)
• Add functions `array_position()` and `array_positions()` to return subscripts of array values (Pavel Stehule)

• Add a point-to-polygon distance operator `<->` (Alexander Korotkov)

• Allow multibyte characters as escapes in `SIMILAR TO` and `SUBSTRING` (Jeff Davis)

  Previously, only a single-byte character was allowed as an escape.

• Add a `width_bucket()` variant that supports any sortable data type and non-uniform bucket widths (Petr Jelínek)

• Add an optional `missing_ok` argument to `pg_read_file()` and related functions (Michael Paquier, Heikki Linnakangas)

• Allow `=>` to specify named parameters in function calls (Pavel Stehule)

  Previously only `:=` could be used. This requires removing the possibility for `=>` to be a user-defined operator. Creation of user-defined `=>` operators has been issuing warnings since PostgreSQL 9.0.

• Add POSIX-compliant rounding for platforms that use PostgreSQL-supplied rounding functions (Pedro Gimeno Fortea)

E.37.3.7.1. System Information Functions and Views

• Add function `pg_get_object_address()` to return OIDs that uniquely identify an object, and function `pg_identify_object_as_address()` to return object information based on OIDs (Álvaro Herrera)

• Loosen security checks for viewing queries in `pg_stat_activity`, executing `pg_cancel_backend()`, and executing `pg_terminate_backend()` (Stephen Frost)

  Previously, only the specific role owning the target session could perform these operations; now membership in that role is sufficient.

• Add `pg_stat_get_snapshot_timestamp()` to output the time stamp of the statistics snapshot (Matt Kelly)

  This represents the last time the snapshot file was written to the file system.

• Add `mxid_age()` to compute multi-xid age (Bruce Momjian)

E.37.3.7.2. Aggregates

• Add `min()/max()` aggregates for `inet/cidr` data types (Haribabu Kommi)

• Use 128-bit integers, where supported, as accumulators for some aggregate functions (Andreas Karlsson)

E.37.3.8. Server-Side Languages

• Improve support for composite types in PL/Python (Ed Behn, Ronan Dunklau)

  This allows PL/Python functions to return arrays of composite types.

• Reduce lossiness of PL/Python floating-point value conversions (Marko Kreen)

• Allow specification of conversion routines between SQL data types and data types of procedural languages (Peter Eisentraut)

  This change adds new commands `CREATE/DROP TRANSFORM`. This also adds optional transformations between the hstore and ltree types to/from PL/Perl and PL/Python.
E.37.3.8.1. PL/pgSQL Server-Side Language

- Improve PL/pgSQL array performance (Tom Lane)
- Add an ASSERT statement in PL/pgSQL (Pavel Stehule)
- Allow more PL/pgSQL keywords to be used as identifiers (Tom Lane)

E.37.3.9. Client Applications

- Move pg_archivecleanup, pg_test_fsync, pg_test_timing, and pg_xlogdump from contrib to src/bin (Peter Eisentraut)
  This should result in these programs being installed by default in most installations.
- Add pg_rewind, which allows re-synchronizing a master server after failback (Heikki Linnakangas)
- Allow pg_receivexlog to manage physical replication slots (Michael Paquier)
  This is controlled via new --create-slot and --drop-slot options.
- Allow pg_receivexlog to synchronously flush WAL to storage using new --synchronous option (Furuya Osamu, Fujii Masao)
  Without this, WAL files are fsync’ed only on close.
- Allow vacuumdb to vacuum in parallel using new --jobs option (Dilip Kumar)
- In vacuumdb, do not prompt for the same password repeatedly when multiple connections are necessary (Haribabu Kommi, Michael Paquier)
- Add --verbose option to reindexdb (Sawada Masahiko)
- Make pg_basebackup use a tablespace mapping file when using tar format, to support symbolic links and file paths of 100+ characters in length on MS Windows (Amit Kapila)
- Add pg_xlogdump option --stats to display summary statistics (Abhijit Menon-Sen)

E.37.3.9.1. psql

- Allow psql to produce AsciiDoc output (Szymon Guz)
- Add an errors mode that displays only failed commands to psql’s ECHO variable (Pavel Stehule)
  This behavior can also be selected with psql’s -b option.
- Provide separate column, header, and border linestyle control in psql’s unicode linestyle (Pavel Stehule)
  Single or double lines are supported; the default is single.
- Add new option %l in psql’s PROMPT variables to display the current multiline statement line number (Sawada Masahiko)
- Add \pset option pager_min_lines to control pager invocation (Andrew Dunstan)
- Improve psql line counting used when deciding to invoke the pager (Andrew Dunstan)
- psql now fails if the file specified by an --output or --log-file switch cannot be written (Tom Lane, Daniel Vérité)
  Previously, it effectively ignored the switch in such cases.
• Add psql tab completion when setting the search_path variable (Jeff Janes)
  Currently only the first schema can be tab-completed.
• Improve psql's tab completion for triggers and rules (Andreas Karlsson)

E.37.3.9.1.1. Backslash Commands

• Add psql \? help sections variables and options (Pavel Stehule)

  \? variables shows psql's special variables and \? options shows the command-line options. \? commands shows the meta-commands, which is the traditional output and remains the default. These help displays can also be obtained with the command-line option --help=section.

• Show tablespace size in psql's \db+ (Fabrízio de Royes Mello)

• Show data type owners in psql's \dT+ (Magnus Hagander)

• Allow psql's \watch to output \timing information (Fujii Masao)
  Also prevent --echo-hidden from echoing \watch queries, since that is generally unwanted.

• Make psql's \sf and \ef commands honor ECHO_HIDDEN (Andrew Dunstan)

• Improve psql tab completion for \set, \unset, and :variable names (Pavel Stehule)

• Allow tab completion of role names in psql \c commands (Ian Barwick)

E.37.3.9.2. pg_dump

• Allow pg_dump to share a snapshot taken by another session using --snapshot (Simon Riggs, Michael Paquier)

  The remote snapshot must have been exported by pg_export_snapshot() or logical replication slot creation. This can be used to share a consistent snapshot across multiple pg_dump processes.

• Support table sizes exceeding 8GB in tar archive format (Tom Lane)

  The POSIX standard for tar format does not allow elements of a tar archive to exceed 8GB, but most modern implementations of tar support an extension that does allow it. Use the extension format when necessary, rather than failing.

• Make pg_dump always print the server and pg_dump versions (Jing Wang)
  Previously, version information was only printed in --verbose mode.

• Remove the long-ignored -i/--ignore-version option from pg_dump, pg_dumpall, and pg_restore (Fujii Masao)

E.37.3.9.3. pg_ctl

• Support multiple pg_ctl -o options, concatenating their values (Bruce Momjian)

• Allow control of pg_ctl's event source logging on MS Windows (MauMau)
  This only controls pg_ctl, not the server, which has separate settings in postgresql.conf.

• If the server's listen address is set to a wildcard value (0.0.0.0 in IPv4 or :: in IPv6), connect via the loopback address rather than trying to use the wildcard address literally (Kondo Yuta)
This fix primarily affects Windows, since on other platforms pg_ctl will prefer to use a Unix-domain socket.

E.37.3.9.4. pg_upgrade

- Move pg_upgrade from contrib to src/bin (Peter Eisentraut)

  In connection with this change, the functionality previously provided by the pg_upgrade_support module has been moved into the core server.

- Support multiple pg_upgrade -c/-o options, concatenating their values (Bruce Momjian)

- Improve database collation comparisons in pg_upgrade (Heikki Linnakangas)

- Remove support for upgrading from 8.3 clusters (Bruce Momjian)

E.37.3.9.5. pgbench

- Move pgbench from contrib to src/bin (Peter Eisentraut)

- Fix calculation of TPS number “excluding connections establishing” (Tatsuo Ishii, Fabien Coelho)

  The overhead for connection establishment was miscalculated whenever the number of pgbench threads was less than the number of client connections. Although this is clearly a bug, we won't back-patch it into pre-9.5 branches since it makes TPS numbers not comparable to previous results.

- Allow counting of pgbench transactions that take over a specified amount of time (Fabien Coelho)

  This is controlled by a new --latency-limit option.

- Allow pgbench to generate Gaussian/exponential distributions using \setrandom (Kondo Mitsumasa, Fabien Coelho)

- Allow pgbench's \set command to handle arithmetic expressions containing more than one operator, and add % (modulo) to the set of operators it supports (Robert Haas, Fabien Coelho)

E.37.3.10. Source Code

- Simplify WAL record format (Heikki Linnakangas)

  This allows external tools to more easily track what blocks are modified.

- Improve the representation of transaction commit and abort WAL records (Andres Freund)

- Add atomic memory operations API (Andres Freund)

- Allow custom path and scan methods (KaiGai Kohei, Tom Lane)

  This allows extensions greater control over the optimizer and executor.

- Allow foreign data wrappers to do post-filter locking (Etsuro Fujita)

- Foreign tables can now take part in INSERT ... ON CONFLICT DO NOTHING queries (Peter Geoghegan, Heikki Linnakangas, Andres Freund)

  Foreign data wrappers must be modified to handle this. INSERT ... ON CONFLICT DO UPDATE is not supported on foreign tables.

- Improve hash_create()’s API for selecting simple-binary-key hash functions (Teodor Sigaev, Tom Lane)
• Improve parallel execution infrastructure (Robert Haas, Amit Kapila, Noah Misch, Rushabh Lathia, Jeevan Chalke)

• Remove Alpha (CPU) and Tru64 (OS) ports (Andres Freund)

• Remove swap-byte-based spinlock implementation for ARMv5 and earlier CPUs (Robert Haas)
  ARMv5’s weak memory ordering made this locking implementation unsafe. Spinlock support is still possible on newer gcc implementations with atomics support.

• Generate an error when excessively long (100+ character) file paths are written to tar files (Peter Eisentraut)
  Tar does not support such overly-long paths.

• Change index operator class for columns pg_seclabel.provider and pg_shseclabel.provider to be text_pattern_ops (Tom Lane)
  This avoids possible problems with these indexes when different databases of a cluster have different default collations.

• Change the spinlock primitives to function as compiler barriers (Robert Haas)

E.37.3.10.1. MS Windows

• Allow higher-precision time stamp resolution on Windows 8, Windows Server 2012, and later Windows systems (Craig Ringer)

• Install shared libraries to bin in MS Windows (Peter Eisentraut, Michael Paquier)

• Install src/test/modules together with contrib on MSVC builds (Michael Paquier)

• Allow configure’s --with-extra-version option to be honored by the MSVC build (Michael Paquier)

• Pass PGFILEDESC into MSVC contrib builds (Michael Paquier)

• Add icons to all MSVC-built binaries and version information to all MS Windows binaries (Noah Misch)
  MinGW already had such icons.

• Add optional-argument support to the internal getopt_long() implementation (Michael Paquier, Andres Freund)
  This is used by the MSVC build.

E.37.3.11. Additional Modules

• Add statistics for minimum, maximum, mean, and standard deviation times to pg_stat_statements
  (Mitsumasa Kondo, Andrew Dunstan)

• Add pgcrypto function pgp_armor_headers() to extract PGP armor headers (Marko Tiikkaja, Heikki Linnakangas)

• Allow empty replacement strings in unaccent (Mohammad Alhashash)
  This is useful in languages where diacritic signs are represented as separate characters.

• Allow multicharacter source strings in unaccent (Tom Lane)
  This could be useful in languages where diacritic signs are represented as separate characters. It also allows more complex unaccent dictionaries.
• Add contrib modules tsm_system_rows and tsm_system_time to allow additional table sampling methods (Petr Jelínek)

• Add GIN index inspection functions to pageinspect (Heikki Linnakangas, Peter Geoghegan, Michael Paquier)

• Add information about buffer pins to pg_buffercache display (Andres Freund)

• Allow pgstattuple to report approximate answers with less overhead using pgstattuple_approx() (Abhijit Menon-Sen)

• Move dummy_seclabel, test_shm_mq, test_parser, and worker_spi from contrib to src/test/modules (Álvaro Herrera)

These modules are only meant for server testing, so they do not need to be built or installed when packaging PostgreSQL.

E.38. Release 9.4.20

Release date: 2018-11-08

This release contains a variety of fixes from 9.4.19. For information about new features in the 9.4 major release, see Section E.58.

E.38.1. Migration to Version 9.4.20

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.18, see Section E.40.

E.38.2. Changes

• Fix corner-case failures in has_foo_privilege() family of functions (Tom Lane)

  Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. has_column_privilege() was additionally capable of crashing on some platforms.

• Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)

• Avoid O(N^3) slowdown in lexer for long strings of + or - characters (Andrew Gierth)

• Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)

• Fix failure of UPDATE/DELETE ... WHERE CURRENT OF ... after rewinding the referenced cursor (Tom Lane)

  A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

• Fix EvalPlanQual to handle conditionally-executed InitPlans properly (Andrew Gierth, Tom Lane)

  This resulted in hard-to-reproduce crashes or wrong answers in concurrent updates, if they contained code such as an uncorrelated sub-SELECT inside a CASE construct.

• Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)
This bug affected full-text-search operations, as well as contrib/ltree and contrib/pg_trgm.

• Ensure that sequences owned by a foreign table are processed by ALTER OWNER on the table (Peter Eisentraut)

  The ownership change should propagate to such sequences as well, but this was missed for foreign tables.

• Fix over-allocation of space for array_out()’s result string (Keiichi Hirobe)

• Fix memory leak in repeated SP-GiST index scans (Tom Lane)

  This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

• Ensure that ApplyLogicalMappingFile() closes the mapping file when done with it (Tomas Vondra)

  Previously, the file descriptor was leaked, eventually resulting in failures during logical decoding.

• Fix logical decoding to handle cases where a mapped catalog table is repeatedly rewritten, e.g. by VACUUM FULL (Andres Freund)

• Prevent starting the server with wal_level set to too low a value to support an existing replication slot (Andres Freund)

• Avoid crash if a utility command causes infinite recursion (Tom Lane)

• When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master (Michael Paquier, Konstantin Knizhnik)

• Randomize the random() seed in bootstrap and standalone backends, and in initdb (Noah Misch)

  The main practical effect of this change is that it avoids a scenario where initdb might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

• Allow DSM allocation to be interrupted (Chris Travers)

• Avoid possible buffer overrun when replaying GIN page recompression from WAL (Alexander Korotkov, Sivasubramanian Ramasubramanian)

• Fix missed fsync of a replication slot’s directory (Konstantin Knizhnik, Michael Paquier)

• Fix unexpected timeouts when using wal_sender_timeout on a slow server (Noah Misch)

• Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)

  This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.

• Don’t run atexit callbacks when servicing SIGQUIT (Heikki Linnakangas)

• Don’t record foreign-server user mappings as members of extensions (Tom Lane)

  If CREATE USER MAPPING is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can’t be extension members, so user mappings shouldn’t be either.

• Make syslogger more robust against failures in opening CSV log files (Tom Lane)
• Fix possible inconsistency in pg_dump's sorting of dissimilar object names (Jacob Champion)

• Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)

This avoids failures due to the new policy of running restores with restrictive search path.

• Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)

pg_upgrade failed to preserve an event trigger's extension-membership status.

• Fix pg_upgrade’s cluster state check to work correctly on a standby server (Bruce Momjian)

• Enforce type cube’s dimension limit in all contrib/cube functions (Andrey Borodin)

Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.

• Fix contrib/unaccent’s unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)

Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.

• Fix build problems on macOS 10.14 (Mojave) (Tom Lane)

Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.

It is now recommended that Perl-related extensions write $(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.

Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.

• Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)

Perl no longer includes the current directory in its search path by default; work around that.

• Support building on Windows with Visual Studio 2015 or Visual Studio 2017 (Michael Paquier, Haribabu Kommi)

• Allow btree comparison functions to return INT_MIN (Tom Lane)

Up to now, we’ve forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of memcmp(), strcmp(), etc, as POSIX doesn’t place any such restriction on those functions. At least some recent versions of memcmp() can return INT_MIN, causing incorrect sort ordering. Hence, we’ve removed this restriction. Callers must now use the INVERTCOMPARE_RESULT() macro if they wish to invert the sort order.

• Fix recursion hazard in shared-invalidation message processing (Tom Lane)

This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.

This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.
• Save and restore SPI's global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)

This prevents possible interference when one SPI-using function calls another.

• Provide ALLOCSET_DEFAULT_SIZES and sibling macros in back branches (Tom Lane)

These macros have existed since 9.6, but there were requests to add them to older branches to allow extensions to rely on them without branch-specific coding.

• Avoid using potentially-under-aligned page buffers (Tom Lane)

Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can't place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.

• Make src/port/snprintf.c follow the C99 standard's definition of snprintf()'s result value (Tom Lane)

On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.

• When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)

This avoids problems with missed floating point overflow checks.

• Fix configure's detection of the result type of strerror_r() (Tom Lane)

The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.

• Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.39. Release 9.4.19

Release date: 2018-08-09

This release contains a variety of fixes from 9.4.18. For information about new features in the 9.4 major release, see Section E.58.

E.39.1. Migration to Version 9.4.19

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.18, see Section E.40.

E.39.2. Changes

• Fix failure to reset libpq's state fully between connection attempts (Tom Lane)

An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general,
an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasickov for reporting this issue. (CVE-2018-10915)

- Ensure that updates to the `relfrozenxid` and `relminmxid` values for “nailed” system catalogs are processed in a timely fashion (Andres Freund)

  Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as `pg_authid`, because the stale cache data could persist into new sessions as well as existing ones.

- Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

  This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

- Avoid emitting a bogus WAL record when recycling an all-zero btree page (Amit Kapila)

  This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

- Improve performance of WAL replay for transactions that drop many relations (Fujii Masao)

  This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

- Improve performance of lock releasing in standby server WAL replay (Thomas Munro)

- Make logical WAL senders report streaming state correctly (Simon Riggs, Sawada Masahiko)

  The code previously mis-detected whether or not it had caught up with the upstream server.

- Fix bugs in snapshot handling during logical decoding, allowing wrong decoding results in rare cases (Arseny Sher, Álvaro Herrera)

- Ensure a table’s cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)

  Previously, the failed index's OID could remain in the list, causing problems later in the same session.

- Fix mishandling of empty uncompressed posting list pages in GIN indexes (Sivasubramanian Ramasubramanian, Alexander Korotkov)

  This could result in an assertion failure after `pg_upgrade` of a pre-9.4 GIN index (9.4 and later will not create such pages).

- Ensure that `VACUUM` will respond to signals within btree page deletion loops (Andres Freund)

  Corrupted btree indexes could result in an infinite loop here, and that previously wasn’t interruptible without forcing a crash.

- Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)

  This resulted in failure to recognize that an index on a composite column could provide the sort order needed for a mergejoin on that column.

- Fix SQL-standard `FETCH FIRST` syntax to allow parameters (`$n`), as the standard expects (Andrew Gierth)
• Fix failure to schema-qualify some object names in getObjectDescription output (Kyotaro Horiguchi, Tom Lane)

Names of collations, conversions, and text search objects were not schema-qualified when they should be.

• Widen COPY FROM's current-line-number counter from 32 to 64 bits (David Rowley)

This avoids two problems with input exceeding 4G lines: COPY FROM WITH HEADER would drop a line every 4G lines, not only the first line, and error reports could show a wrong line number.

• Add a string freeing function to ecpg's pgtypes library, so that cross-module memory management problems can be avoided on Windows (Takayuki Tsunakawa)

On Windows, crashes can ensue if the free call for a given chunk of memory is not made from the same DLL that malloc'ed the memory. The pgtypes library sometimes returns strings that it expects the caller to free, making it impossible to follow this rule. Add a PGTYPESchar_free() function that just wraps free, allowing applications to follow this rule.

• Fix ecpg's support for long long variables on Windows, as well as other platforms that declare strtoll/strtoull nonstandardly or not at all (Dang Minh Huong, Tom Lane)

• Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in the semantics of a statement intra-session (Tom Lane)

This error led to assertion failures, or in rare cases, failure to enforce the INTO STRICT option as expected.

• Fix password prompting in client programs so that echo is properly disabled on Windows when stdin is not the terminal (Matthew Stickney)

• Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)

The previous fix for quoting of search_path and other list-valued variables in pg_dump output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

• Fix pg_dump's failure to dump REPLICA IDENTITY properties for constraint indexes (Tom Lane)

Manually created unique indexes were properly marked, but not those created by declaring UNIQUE or PRIMARY KEY constraints.

• Make pg_upgrade check that the old server was shut down cleanly (Bruce Momjian)

The previous check could be fooled by an immediate-mode shutdown.

• Fix crash in contrib/ltree's lca() function when the input array is empty (Pierre Ducroquet)

• Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

• Rearrange makefiles to ensure that programs link to freshly-built libraries (such as libpq.so) rather than ones that might exist in the system library directories (Tom Lane)

This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

• Update time zone data files to tzdata release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.

This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbrev-
viations in use; the only known effect is that the is_dst column in the pg_timezone_names view will now be true in winter and false in summer in these cases.

E.40. Release 9.4.18

Release date: 2018-05-10

This release contains a variety of fixes from 9.4.17. For information about new features in the 9.4 major release, see Section E.58.

E.40.1. Migration to Version 9.4.18

A dump/restore is not required for those running 9.4.X.

However, if the function marking mistakes mentioned in the first changelog entry below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 9.4.17, see Section E.41.

E.40.2. Changes

- Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

  The functions query_to_xml, cursor_to_xml, cursor_to_xmlschema, query_to_xmlschema, and query_to_xml_and_xmlschema should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions’ pg_proc entries, for example ALTER FUNCTION pg_catalog.query_to_xml(text, boolean, boolean, text) VOLATILE. (Note that that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

- Avoid re-using TOAST value OIDs that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

  Once the OID counter has wrapped around, it's possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value nnnnn” errors, which would persist until the dead entry was removed by VACUUM. Fix by not selecting such OIDs when creating a new TOAST entry.

- Change ANALYZE's algorithm for updating pg_class.reltuples (David Gould)

  Previously, pages not actually scanned by ANALYZE were assumed to retain their old tuple density. In a large table where ANALYZE samples only a small fraction of the pages, this meant that the overall tuple density estimate could not change very much, so that reltuples would change nearly proportionally to changes in the table’s physical size (relpages) regardless of what was actually happening in the table. This has been observed to result in reltuples becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that ANALYZE's sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

- Avoid deadlocks in concurrent CREATE INDEX CONCURRENTLY commands that are run under SERIALIZABLE or REPEATABLE READ transaction isolation (Tom Lane)

- Fix possible slow execution of REFRESH MATERIALIZED VIEW CONCURRENTLY (Thomas Munro)
• Fix **UPDATE/DELETE** ... **WHERE CURRENT OF** to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)

• Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)
  This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.

• Fix misoptimization of **CHECK** constraints having provably-NULL subclauses of top-level **AND/OR** conditions (Tom Lane, Dean Rasheed)
  This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.

• Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)

• Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)

• Fix overly strict sanity check in **heap_prepare_freeze_tuple** (Álvaro Herrera)
  This could result in incorrect “cannot freeze committed xmax” failures in databases that have been **pg_upgrade**d from 9.2 or earlier.

• Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)

• Reduce locking during autovacuum worker scheduling (Jeff Janes)
  The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.

• Ensure client hostname is copied while copying **pg_stat_activity** data to local memory (Edmund Horner)
  Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the client hostname column to change unexpectedly if any existing session disconnected.

• Fix incorrect processing of multiple compound affixes in **ispell** dictionaries (Arthur Zakirov)

• Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on text columns (Tom Lane)
  Such searches would return the wrong set of rows in most non-C locales.

• Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)
  Previously, the tuple count was reported to be the same as that of the underlying table, which is wrong if the index is partial.

• Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)
  Previously it reported the estimated number of heap tuples, which might be inaccurate, and is certainly wrong if the index is partial.

• Fix a corner case where a streaming standby gets stuck at a WAL continuation record (Kyotaro Horiguchi)

• In logical decoding, avoid possible double processing of WAL data when a walsender restarts (Craig Ringer)

• Allow **scalarltsel** and **scalargtsel** to be used on non-core datatypes (Tomas Vondra)
• Reduce libpq's memory consumption when a server error is reported after a large amount of query output has been collected (Tom Lane)

Discard the previous output before, not after, processing the error message. On some platforms, notably Linux, this can make a difference in the application's subsequent memory footprint.

• Fix double-free crashes in ecpg (Patrick Krecker, Jeevan Ladhe)

• Fix ecpg to handle long long int variables correctly in MSVC builds (Michael Meskes, Andrew Gierth)

• Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)

The local_preload_libraries, session_preload_libraries, shared_preload_libraries, and temp_tablespaces variables were not correctly quoted in pg_dump output. This would cause problems if settings for these variables appeared in CREATE FUNCTION ... SET or ALTER DATABASE/ROLE ... SET clauses.

• Fix pg_recvlogical to not fail against pre-v10 PostgreSQL servers (Michael Paquier)

A previous fix caused pg_recvlogical to issue a command regardless of server version, but it should only be issued to v10 and later servers.

• Fix overflow handling in PL/pgSQL integer FOR loops (Tom Lane)

The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

• Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisentraut)

• Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)

• Rename internal b64_encode and b64_decode functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)

• Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)

This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.

• Update time zone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

E.41. Release 9.4.17

Release date: 2018-03-01

This release contains a variety of fixes from 9.4.16. For information about new features in the 9.4 major release, see Section E.58.

E.41.1. Migration to Version 9.4.17

A dump/restore is not required for those running 9.4.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that
you read the documentation changes described in the first changelog entry below, and take suitable
steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index
expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After
upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 9.4.13, see Section E.45.

E.41.2. Changes

• Document how to configure installations and applications to guard against search-path-dependent
trojan-horse attacks from other users (Noah Misch)

Using a search_path setting that includes any schemas writable by a hostile user enables that
user to capture control of queries and then run arbitrary SQL code with the permissions of the at-
tacked user. While it is possible to write queries that are proof against such hijacking, it is notation-
ally tedious, and it's very easy to overlook holes. Therefore, we now recommend configurations
in which no untrusted schemas appear in one's search path. Relevant documentation appears in
Section 5.8.6 (for database administrators and users), Section 34.1 (for application authors), Sec-
tion 38.16.1 (for extension authors), and CREATE FUNCTION (for authors of SECURITY DE-
FINER functions). (CVE-2018-1058)

• Avoid use of insecure search_path settings in pg_dump and other client programs (Noah Misch, Tom Lane)

pg_dump, pg_upgrade, vacuumdb and other PostgreSQL-provided applications were themselves
vulnerable to the type of hijacking described in the previous changelog entry; since these applica-
tions are commonly run by superusers, they present particularly attractive targets. To make them
secure whether or not the installation as a whole has been secured, modify them to include only
the pg_catalog schema in their search_path settings. Autovacuum worker processes now
do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example,
user-provided functions in index expressions — the tighter search_path may result in errors,
which will need to be corrected by adjusting those user-provided functions to not assume anything
about what search path they are invoked under. That has always been good practice, but now it will
be necessary for correct behavior. (CVE-2018-1058)

• Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom
Lane)

If a CTE (WITH clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck
due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

• Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do
not match mergeclauses” planner errors in corner cases.

• Repair pg_upgrade's failure to preserve relfrozenxid for materialized views (Tom Lane, And-
dres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting
as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The
problem would be more likely to occur in seldom-refreshed materialized views, or ones that were
maintained only with REFRESH MATERIALIZED VIEW CONCURRENTLY.

If such corruption is observed, it can be repaired by refreshing the materialized view (without CON-
CURRENTLY).
Release Notes

- Fix incorrect reporting of PL/Python function names in error CONTEXT stacks (Tom Lane)

  An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function's name twice, rather than the expected results. Also, an error in a nested PL/Python DO block could result in a null pointer dereference crash on some platforms.

- Allow contrib/auto_explain's log_min_duration setting to range up to INT_MAX, or about 24 days instead of 35 minutes (Tom Lane)

E.42. Release 9.4.16

Release date: 2018-02-08

This release contains a variety of fixes from 9.4.15. For information about new features in the 9.4 major release, see Section E.58.

E.42.1. Migration to Version 9.4.16

A dump/restore is not required for those running 9.4.X. However, if you are upgrading from a version earlier than 9.4.13, see Section E.45.

E.42.2. Changes

- Ensure that all temporary files made by pg_upgrade are non-world-readable (Tom Lane, Noah Misch)

  pg_upgrade normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing pg_dumpall -g output would be group- or world-readable, or even writable, if the user's umask setting allows. In typical usage on multi-user machines, the umask and/or the working directory's permissions would be tight enough to prevent problems; but there may be people using pg_upgrade in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

- Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

  In some cases VACUUM would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

- Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)

  These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.

- Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by ALTER TABLE NO INHERIT (Tom Lane)

- Fix spurious deadlock failures when multiple sessions are running CREATE INDEX CONCURRENTLY (Jeff Janes)

- Repair failure with correlated sub-SELECT inside VALUES inside a LATERAL subquery (Tom Lane)

- Fix “could not devise a query plan for the given query” planner failure for some cases involving nested UNION ALL inside a lateral subquery (Tom Lane)

- Fix logical decoding to correctly clean up disk files for crashed transactions (Atsushi Torikoshi)
Logical decoding may spill WAL records to disk for transactions generating many WAL records. Normally these files are cleaned up after the transaction's commit or abort record arrives; but if no such record is ever seen, the removal code misbehaved.

- Fix walsender timeout failure and failure to respond to interrupts when processing a large transaction (Petr Jelinek)
- Fix has_sequence_privilege() to support WITH GRANT OPTION tests, as other privilege-testing functions do (Joe Conway)
- In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)

We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don't promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only xpath() and related functions; other XML code paths already acted this way.

- Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)

Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.

- Cope with failure to start a parallel worker process (Amit Kapila, Robert Haas)

Parallel query previously tended to hang indefinitely if a worker could not be started, as the result of fork() failure or other low-probability problems.

- Prevent stack-overflow crashes when planning extremely deeply nested set operations (UNION/INTERSECT/EXCEPT) (Tom Lane)

- Fix null-pointer crashes for some types of LDAP URLs appearing in pg_hba.conf (Thomas Munro)

- Fix sample INSTR() functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)

These functions are stated to be Oracle® compatible, but they weren't exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.

The sample code has been adjusted to match Oracle's behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

- Fix pg_dump to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with DATABASE so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with LARGE OBJECT. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.
Note that this change has user-visible results in the output of `pg_restore --list`.

- In ecpg, detect indicator arrays that do not have the correct length and report an error (David Rader)
- Avoid triggering a libc assertion in `contrib/hstore`, due to use of `memcpy()` with equal source and destination pointers (Tomas Vondra)
- Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

The scripts in `contrib/start-scripts/osx` use infrastructure that's been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory `contrib/start-scripts/macos` containing scripts that use the newer launchd infrastructure.

- Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dunstan)
- Support linking to MinGW-built versions of libperl (Noah Misch)
  This allows building PL/Perl with some common Perl distributions for Windows.
- Fix MSVC build to test whether 32-bit libperl needs `~D_USE_32BIT_TIME_T` (Noah Misch)
  Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.
- On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)
  This may allow collection of a core dump for some early-startup failures that did not produce a dump before.
- On Windows, avoid encoding-conversion-related crashes when emitting messages very early in postmaster startup (Takayuki Tsunakawa)
- Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)
- Add support for spinlocks on Motorola 88K (David Carlier)
- Update time zone data files to `tzdata` release 2018c for DST law changes in Brazil, Sao Tome and Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The `US/Pacific-New` zone has been removed (it was only an alias for `America/Los_Angeles` anyway).

### E.43. Release 9.4.15

**Release date:** 2017-11-09

This release contains a variety of fixes from 9.4.14. For information about new features in the 9.4 major release, see Section E.58.

### E.43.1. Migration to Version 9.4.15

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.13, see Section E.45.

### E.43.2. Changes

- Fix crash due to rowtype mismatch in `json{b}_populate_recordset()` (Michael Paquier, Tom Lane)
These functions used the result rowtype specified in the FROM ... AS clause without checking that it matched the actual rowtype of the supplied tuple value. If it didn't, that would usually result in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)

• Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)

Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making $PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

• Fix crash when logical decoding is invoked from a SPI-using function, in particular any function written in a PL language (Tom Lane)

• Fix json_build_array(), json_build_object(), and their jsonb equivalents to handle explicit VARIADIC arguments correctly (Michael Paquier)

• Properly reject attempts to convert infinite float values to type numeric (Tom Lane, KaiGai Kohei)

Previously the behavior was platform-dependent.

• Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)

• Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes (Tom Lane)

Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

• Correctly detect hashability of range data types (Tom Lane)

The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range's subtype has hash support. This does not affect any of the built-in range types, since they're all hashable anyway.

• Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)

If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.

• Avoid SIGBUS crash on Linux when a DSM memory request exceeds the space available in tmpfs (Thomas Munro)

• Prevent low-probability crash in processing of nested trigger firings (Tom Lane)

• Allow COPY's FREEZE option to work when the transaction isolation level is REPEATABLE READ or higher (Noah Misch)

This case was unintentionally broken by a previous bug fix.

• Correctly restore the umask setting when file creation fails in COPY or lo_export() (Peter Eisentraut)

• Give a better error message for duplicate column names in ANALYZE (Nathan Bossart)

• Fix mis-parsing of the last line in a non-newline-terminated pg_hba.conf file (Tom Lane)
• Fix libpq to not require user's home directory to exist (Tom Lane)

In v10, failure to find the home directory while trying to read ~/.pgpass was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading ~/.pg_service.conf, though this was less obvious since that file is not sought unless a service name is specified.

• Fix libpq to guard against integer overflow in the row count of a PGresult (Michael Paquier)

• Fix ecpg's handling of out-of-scope cursor declarations with pointer or array variables (Michael Meskes)

• In ecpglib, correctly handle backslashes in string literals depending on whether standard_conforming_strings is set (Tsunakawa Takayuki)

• Make ecpglib's Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengu, Michael Meskes)

• Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

This fixes various issues: the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no posixrules file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.

• Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

E.44. Release 9.4.14

Release date: 2017-08-31

This release contains a small number of fixes from 9.4.13. For information about new features in the 9.4 major release, see Section E.58.

E.44.1. Migration to Version 9.4.14

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.13, see Section E.45.

E.44.2. Changes

• Fix failure of walsender processes to respond to shutdown signals (Marco Nenciarini)

A missed flag update resulted in walsenders continuing to run as long as they had a standby server connected, preventing primary-server shutdown unless immediate shutdown mode is used.

• Show foreign tables in information_schema.table_privileges view (Peter Eisentraut)

All other relevant information_schema views include foreign tables, but this one ignored them.

Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:

```
SET search_path TO information_schema;
CREATE OR REPLACE VIEW table_privileges AS
```
SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
CAST(grantee.rolname AS sql_identifier) AS grantee,
CAST(current_database() AS sql_identifier) AS table_catalog,
CAST(nc.nspname AS sql_identifier) AS table_schema,
CAST(c.relname AS sql_identifier) AS table_name,
CAST(c.prtype AS character_data) AS privilege_type,
CAST(
  CASE WHEN
      -- object owner always has grant options
      pg_has_role(grantee.oid, c.relowner, 'USAGE')
      OR c.grantable
      THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE
  'NO' END AS yes_or_no) AS with_hierarchy
FROM (SELECT oid, relname, relnamespace, relkind, relowner,
  (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class ) AS c (oid, relname, relnamespace, relkind, relowner, grantor, grantee, prtype, grantable),
pg_namespace nc,
pag_authid u_grantor,
{ SELECT oid, rolname FROM pag_authid
  UNION ALL
  SELECT 0::oid, 'PUBLIC'
 } AS grantee (oid, rolname)
WHERE c.relnamespace = nc.oid
  AND c.relkind IN ('r', 'v', 'f')
  AND c.grantee = grantee.oid
  AND c.grantor = u_grantor.oid
  AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE',
    'DELETE', 'TRUNCATE', 'REFERENCES', 'TRIGGER')
  AND (pg_has_role(u_grantor.oid, 'USAGE')
    OR pg_has_role(grantee.oid, 'USAGE')
    OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

• Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)
  
  This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about “cannot drop active portal”.

• Remove assertion that could trigger during a fatal exit (Tom Lane)

• Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

  Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant values that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.
• Fix crash in pg_restore when using parallel mode and using a list file to select a subset of items to restore (Fabrizio de Royes Mello)

• Change ecpg's parser to allow RETURNING clauses without attached C variables (Michael Meskes)

  This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

• Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

  This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that _USE_32BIT_TIME_T is no longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

---

**E.45. Release 9.4.13**

**Release date:** 2017-08-10

This release contains a variety of fixes from 9.4.12. For information about new features in the 9.4 major release, see Section E.58.

**E.45.1. Migration to Version 9.4.13**

A dump/restore is not required for those running 9.4.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.4.12, see Section E.46.

**E.45.2. Changes**

• Further restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Noah Misch)

  The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have USAGE permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since information_schema.user_mapping_options does not show the options in such cases, pg_user_mappings should not either. (CVE-2017-7547)

  By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, you will need to do the following:

  1. Restart the postmaster after adding allow_system_table_mods = true to postgresql.conf. (In versions supporting ALTER SYSTEM, you can use that to make the configuration change, but you'll still need a restart.)

  2. In each database of the cluster, run the following commands as superuser:

  ```sql
  SET search_path = pg_catalog;
  CREATE OR REPLACE VIEW pg_user_mappings AS
  SELECT
      U.oid       AS umid,
      S.oid       AS srvid,
      S.srvname   AS srvname,
  FROM
      pg_user_mappings U,
      pg_service_addresses S
  WHERE
      U.umid = S.umid
      AND S.srvname = S.srvname
      AND S.srvname = S.srvname;
  ```
U.umuser    AS umuser,
CASE WHEN U.umuser = 0 THEN
  'public'
ELSE
  A.rolname
END AS usename,
CASE WHEN (U.umuser <> 0 AND A.rolname = current_user
  AND (pg_has_role(S.srvowner, 'USAGE')
       OR has_server_privilege(S.oid,
          'USAGE')))
  OR (U.umuser = 0 AND
      pg_has_role(S.srvowner, 'USAGE'))
  OR (SELECT rolsuper FROM pg_authid WHERE
       rolname = current_user)
  THEN U.umoptions
  ELSE NULL END AS umoptions
FROM pg_user_mapping U
  LEFT JOIN pg_authid A ON (A.oid = U.umuser)
  JOIN pg_foreign_server S ON (U.umserver = S.oid);

3. Do not forget to include the template0 and template1 databases, or the vulnerability will still exist in databases you create later. To fix template0, you'll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;

and then after fixing template0, undo that with

ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;

In prior versions, instead use

UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';

4. Finally, remove the allow_system_tableMods configuration setting, and again restart the postmaster.

• Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)

libpq ignores empty password specifications, and does not transmit them to the server. So, if a user's password has been set to the empty string, it's impossible to log in with that password via psql or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, md5, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

• Make lo_put() check for UPDATE privilege on the target large object (Tom Lane, Michael Paquier)

lo_put() should surely require the same permissions as lowrite(), but the check was missing, allowing any user to change the data in a large object. (CVE-2017-7548)

• Fix concurrent locking of tuple update chains (Álvaro Herrera)
If several sessions concurrently lock a tuple update chain with nonconflicting lock modes using an old snapshot, and they all succeed, it was possible for some of them to nonetheless fail (and conclude there is no live tuple version) due to a race condition. This had consequences such as foreign-key checks failing to see a tuple that definitely exists but is being updated concurrently.

- Fix potential data corruption when freezing a tuple whose XMAX is a multixact with exactly one still-interesting member (Teodor Sigaev)
- Avoid integer overflow and ensuing crash when sorting more than one billion tuples in-memory (Sergey Koposov)
- On Windows, retry process creation if we fail to reserve the address range for our shared memory in the new process (Tom Lane, Amit Kapila)
  
  This is expected to fix infrequent child-process-launch failures that are probably due to interference from antivirus products.
- Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas Munro, Tom Lane)
- Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael Paquier)
- Prevent sending SSL session tickets to clients (Tom Lane)
  
  This fix prevents reconnection failures with ticket-aware client-side SSL code.
- Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)
- Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and immediate restart (Tom Lane)
  
  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively ignored.
- Ensure that the statistics collector's receive buffer size is at least 100KB (Tom Lane)
  
  This reduces the risk of dropped statistics data on older platforms whose default receive buffer size is less than that.
- Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes an XLOG_SWITCH WAL record (Andres Freund)
- Fix walsender to exit promptly when client requests shutdown (Tom Lane)
- Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)
- Prevent walsender-triggered panics during shutdown checkpoints (Andres Freund, Michael Paquier)
- Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom Lane)
- Fix logical decoding failure with very wide tuples (Andres Freund)
  
  Logical decoding crashed on tuples that are wider than 64KB (after compression, but with all data in-line). The case arises only when REPLICA IDENTITY FULL is enabled for a table containing such tuples.
- Fix leakage of small subtransactions spilled to disk during logical decoding (Andres Freund)
This resulted in temporary files consuming excessive disk space.

- Reduce the work needed to build snapshots during creation of logical-decoding slots (Andres Freund, Petr Jelinek)

The previous algorithm was infeasibly expensive on a server with a lot of open transactions.

- Fix race condition that could indefinitely delay creation of logical-decoding slots (Andres Freund, Petr Jelinek)

- Reduce overhead in processing syscache invalidation events (Tom Lane)

  This is particularly helpful for logical decoding, which triggers frequent cache invalidation.

- Fix cases where an INSERT or UPDATE assigns to more than one element of a column that is of domain-over-array type (Tom Lane)

- Allow window functions to be used in sub-SELECTs that are within the arguments of an aggregate function (Tom Lane)

- Move autogenerated array types out of the way during ALTER ... RENAME (Vik Fearing)

  Previously, we would rename a conflicting autogenerated array type out of the way during CREATE; this fix extends that behavior to renaming operations.

- Ensure that ALTER USER ... SET accepts all the syntax variants that ALTER ROLE ... SET does (Peter Eisentraut)

- Properly update dependency info when changing a datatype I/O function's argument or return type from opaque to the correct type (Heikki Linnakangas)

  CREATE TYPE updates I/O functions declared in this long-obsolete style, but it forgot to record a dependency on the type, allowing a subsequent DROP TYPE to leave broken function definitions behind.

- Reduce memory usage when ANALYZE processes a tsvector column (Heikki Linnakangas)

- Fix unnecessary precision loss and sloppy rounding when multiplying or dividing money values by integers or floats (Tom Lane)

- Tighten checks for whitespace in functions that parse identifiers, such as regprocedurein() (Tom Lane)

  Depending on the prevailing locale, these functions could misinterpret fragments of multibyte characters as whitespace.

- Use relevant #define symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)

  This avoids portability problems, typically manifesting as a “handshake” mismatch during library load, when working with recent Perl versions.

- In libpq, reset GSS/SASL and SSPI authentication state properly after a failed connection attempt (Michael Paquier)

  Failure to do this meant that when falling back from SSL to non-SSL connections, a GSS/SASL failure in the SSL attempt would always cause the non-SSL attempt to fail. SSPI did not fail, but it leaked memory.

- In psql, fix failure when COPY FROM STDIN is ended with a keyboard EOF signal and then another COPY FROM STDIN is attempted (Thomas Munro)
This misbehavior was observed on BSD-derived platforms (including macOS), but not on most others.

- Fix pg_dump and pg_restore to emit REFRESH MATERIALIZED VIEW commands last (Tom Lane)

This prevents errors during dump/restore when a materialized view refers to tables owned by a different user.

- Improve pg_dump/pg_restore's reporting of error conditions originating in zlib (Vladimir Kuruschikov, Álvaro Herrera)

- Fix pg_dump with the --clean option to drop event triggers as expected (Tom Lane)

It also now correctly assigns ownership of event triggers; before, they were restored as being owned by the superuser running the restore script.

- Fix pg_dump to not emit invalid SQL for an empty operator class (Daniel Gustafsson)

- Fix pg_dump output to stdout on Windows (Kuntal Ghosh)

A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_get_ruledef() to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)

In some corner cases, pg_dump relies on pg_get_ruledef() to dump views, so that this error could result in dump/reload failures.

- Fix dumping of outer joins with empty constraints, such as the result of a NATURAL LEFT JOIN with no common columns (Tom Lane)

- Fix dumping of function expressions in the FROM clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)

- Fix pg_basebackup output to stdout on Windows (Haribabu Kommi)

A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_upgrade to ensure that the ending WAL record does not have wal_level = minimum (Bruce Momjian)

This condition could prevent upgraded standby servers from reconnecting.

- In postgres_fdw, re-establish connections to remote servers after ALTER SERVER or ALTER USER MAPPING commands (Kyotaro Horiguchi)

This ensures that option changes affecting connection parameters will be applied promptly.

- In postgres_fdw, allow cancellation of remote transaction control commands (Robert Haas, Rafia Sabih)

This change allows us to quickly escape a wait for an unresponsive remote server in many more cases than previously.

- Increase MAX_SYSCACHE_CALLBACKS to provide more room for extensions (Tom Lane)

- Always use -fPIC, not -fpic, when building shared libraries with gcc (Tom Lane)
This supports larger extension libraries on platforms where it makes a difference.

- Fix unescaped-braces issue in our build scripts for Microsoft MSVC, to avoid a warning or error from recent Perl versions (Andrew Dunstan)

- In MSVC builds, handle the case where the OpenSSL library is not within a VC subdirectory (Andrew Dunstan)

- In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)
  This fixes a former need to move things around in standard Windows installations of libxml2.

- In MSVC builds, recognize a Tcl library that is named tcl86.lib (Noah Misch)

- In MSVC builds, honor PROVE_FLAGS settings on vcregress.pl's command line (Andrew Dunstan)

### E.46. Release 9.4.12

**Release date:** 2017-05-11

This release contains a variety of fixes from 9.4.11. For information about new features in the 9.4 major release, see Section E.58.

#### E.46.1. Migration to Version 9.4.12

A dump/restore is not required for those running 9.4.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are using third-party replication tools that depend on “logical decoding”, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.4.11, see Section E.47.

#### E.46.2. Changes

- Restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)

  The previous coding allowed the owner of a foreign server object, or anyone he has granted server USAGE permission to, to see the options for all user mappings associated with that server. This might well include passwords for other users. Adjust the view definition to match the behavior of information_schema.user_mapping_options, namely that these options are visible to the user being mapped, or if the mapping is for PUBLIC and the current user is the server owner, or if the current user is a superuser. (CVE-2017-7486)

  By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, follow the corrected procedure shown in the changelog entry for CVE-2017-7547, in Section E.45.

- Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

  Some selectivity estimation functions in the planner will apply user-defined operators to values obtained from pg_statistic, such as most common values and histogram entries. This occurs before table permissions are checked, so a nefarious user could exploit the behavior to obtain these values for table columns he does not have permission to read. To fix, fall back to a default estimate if the operator's implementation function is not certified leak-proof and the calling user does not
have permission to read the table column whose statistics are needed. At least one of these criteria is satisfied in most cases in practice. (CVE-2017-7484)

• Restore libpq's recognition of the PGREQUIRESSL environment variable (Daniel Gustafsson)

Processing of this environment variable was unintentionally dropped in PostgreSQL 9.3, but its documentation remained. This creates a security hazard, since users might be relying on the environment variable to force SSL-encrypted connections, but that would no longer be guaranteed. Restore handling of the variable, but give it lower priority than PGSSLMODE, to avoid breaking configurations that work correctly with post-9.3 code. (CVE-2017-7485)

• Fix possibly-invalid initial snapshot during logical decoding (Petr Jelinek, Andres Freund)

The initial snapshot created for a logical decoding replication slot was potentially incorrect. This could cause third-party tools that use logical decoding to copy incomplete/inconsistent initial data. This was more likely to happen if the source server was busy at the time of slot creation, or if another logical slot already existed.

If you are using a replication tool that depends on logical decoding, and it should have copied a nonempty data set at the start of replication, it is advisable to recreate the replica after installing this update, or to verify its contents against the source server.

• Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)

This could result in an unlogged index being set to an invalid state after a crash and restart. Such a problem would persist until the index was dropped and rebuilt.

• Fix incorrect reconstruction of pg_subtrans entries when a standby server replays a prepared but uncommitted two-phase transaction (Tom Lane)

In most cases this turned out to have no visible ill effects, but in corner cases it could result in circular references in pg_subtrans, potentially causing infinite loops in queries that examine rows modified by the two-phase transaction.

• Avoid possible crash in walsender due to failure to initialize a string buffer (Stas Kelvich, Fujii Masao)

• Fix postmaster's handling of fork() failure for a background worker process (Tom Lane)

Previously, the postmaster updated portions of its state as though the process had been launched successfully, resulting in subsequent confusion.

• Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL (Julien Rouhaud, Tom Lane)

Due to lack of a cache flush step between commands in an extension script file, non-utility queries might not see the effects of an immediately preceding catalog change, such as ALTER TABLE ... RENAME.

• Skip tablespace privilege checks when ALTER TABLE ... ALTER COLUMN TYPE rebuilds an existing index (Noah Misch)

The command failed if the calling user did not currently have CREATE privilege for the tablespace containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be rebuilt where it is.

• Fix ALTER TABLE ... VALIDATE CONSTRAINT to not recurse to child tables when the constraint is marked NO INHERIT (Amit Langote)

This fix prevents unwanted “constraint does not exist” failures when no matching constraint is present in the child tables.
• Fix VACUUM to account properly for pages that could not be scanned due to conflicting page pins (Andrew Gierth)

This tended to lead to underestimation of the number of tuples in the table. In the worst case of a small heavily-contended table, VACUUM could incorrectly report that the table contained no tuples, leading to very bad planning choices.

• Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests (Tom Lane, Thomas Munro)

• Fix integer-overflow problems in interval comparison (Kyotaro Horiguchi, Tom Lane)

The comparison operators for type interval could yield wrong answers for intervals larger than about 296000 years. Indexes on columns containing such large values should be reindexed, since they may be corrupt.

• Fix cursor_to_xml() to produce valid output with tableforest = false (Thomas Munro, Peter Eisentraut)

Previously it failed to produce a wrapping <table> element.

• Fix roundoff problems in float8_timestamptz() and make_interval() (Tom Lane)

These functions truncated, rather than rounded, when converting a floating-point value to integer microseconds; that could cause unexpectedly off-by-one results.

• Improve performance of pg_timezone_names view (Tom Lane, David Rowley)

• Reduce memory management overhead for contexts containing many large blocks (Tom Lane)

• Fix sloppy handling of corner-case errors from lseek() and close() (Tom Lane)

Neither of these system calls are likely to fail in typical situations, but if they did, fd.c could get quite confused.

• Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)

This could result in attempting to write to the event log when that isn't accessible, so that no logging happens at all.

• Fix ecpg to support COMMIT PREPARED and ROLLBACK PREPARED (Masahiko Sawada)

• Fix a double-free error when processing dollar-quoted string literals in ecpg (Michael Meskes)

• In pg_dump, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)

In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.

• Avoid emitting an invalid list file in pg_restore -l when SQL object names contain newlines (Tom Lane)

Replace newlines by spaces, which is sufficient to make the output valid for pg_restore -L's purposes.

• Fix pg_upgrade to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)

Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.
• Improve error handling in contrib/adminpack's pg_file_write() function (Noah Misch)

Notably, it failed to detect errors reported by fclose().

• In contrib/dblink, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

• Fix contrib/pg_trgm's extraction of trigrams from regular expressions (Tom Lane)

In some cases it would produce a broken data structure that could never match anything, leading to GIN or GiST indexscans that use a trigram index not finding any matches to the regular expression.

• In contrib/postgres_fdw, transmit query cancellation requests to the remote server (Michael Paquier, Etsuro Fujita)

Previously, a local query cancellation request did not cause an already-sent remote query to terminate early. This is a back-patch of work originally done for 9.6.

• Support OpenSSL 1.1.0 (Heikki Linnakangas, Andreas Karlsson, Tom Lane)

This is a back-patch of work previously done in newer branches; it's needed since many platforms are adopting newer OpenSSL versions.

• Support Tcl 8.6 in MSVC builds (Álvaro Herrera)

• Sync our copy of the timezone library with IANA release tzcode2017b (Tom Lane)

This fixes a bug affecting some DST transitions in January 2038.

• Update time zone data files to tzdata release 2017b for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

• Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds (David Rowley)

The Microsoft MSVC build scripts neglected to install the posixrules file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the US/Eastern zone. If you don't want that, remove the posixrules file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.

E.47. Release 9.4.11

Release date: 2017-02-09

This release contains a variety of fixes from 9.4.10. For information about new features in the 9.4 major release, see Section E.58.
E.47.1. Migration to Version 9.4.11

A dump/restore is not required for those running 9.4.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.4.10, see Section E.48.

E.47.2. Changes

• Fix a race condition that could cause indexes built with \texttt{CREATE INDEX CONCURRENTLY} to be corrupt (Pavan Deolasee, Tom Lane)

If \texttt{CREATE INDEX CONCURRENTLY} was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the \texttt{CREATE INDEX} command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

• Ensure that the special snapshot used for catalog scans is not invalidated by premature data pruning (Tom Lane)

Backends failed to account for this snapshot when advertising their oldest xmin, potentially allowing concurrent vacuuming operations to remove data that was still needed. This led to transient failures along the lines of “cache lookup failed for relation 1255”.

• Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)

Previously, this was skipped when \texttt{wal_level = minimal}, but actually it's necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

• Reduce interlocking on standby servers during the replay of btree index vacuuming operations (Simon Riggs)

This change avoids substantial replication delays that sometimes occurred while replaying such operations.

• If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)

• Ensure that hot standby feedback works correctly when it's enabled at standby server start (Ants Asmaa, Craig Ringer)

• Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

• Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)

This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

• Fix check for when an extension member object can be dropped (Tom Lane)

Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

• Make sure \texttt{ALTER TABLE} preserves index tablespace assignments when rebuilding indexes (Tom Lane, Michael Paquier)

Previously, non-default settings of default_tablespace could result in broken indexes.

• Fix incorrect updating of trigger function properties when changing a foreign-key constraint's deferrability properties with \texttt{ALTER TABLE \ldots ALTER CONSTRAINT} (Tom Lane)
This led to odd failures during subsequent exercise of the foreign key, as the triggers were fired at the wrong times.

- Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)

  This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.

- Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via ALTER TABLE ... INHERIT (Amit Langote)

  The OID column should be treated the same as regular user columns in this case, but it wasn't, leading to odd behavior in later inheritance changes.

- Fix CREATE OR REPLACE VIEW to update the view query before attempting to apply the new view options (Dean Rasheed)

  Previously the command would fail if the new options were inconsistent with the old view definition.

- Report correct object identity during ALTER TEXT SEARCH CONFIGURATION (Artur Zakirov)

  The wrong catalog OID was reported to extensions such as logical decoding.

- Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

  When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

- Prevent multicolumn expansion of foo.* in an UPDATE source expression (Tom Lane)

  This led to “UPDATE target count mismatch --- internal error”. Now the syntax is understood as a whole-row variable, as it would be in other contexts.

- Ensure that column typmods are determined accurately for multi-row VALUES constructs (Tom Lane)

  This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a varchar value) but later values don't share the same limit.

- Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

  Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (U+"...") or Unicode identifier (U+"...").

- Ensure that a purely negative text search query, such as !foo, matches empty ts vectors (Tom Dunstan)

  Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

- Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

- Fix performance problems in ts_rewrite() (Tom Lane)

- Fix ts_rewrite()'s handling of nested NOT operators (Tom Lane)
• Fix `array_fill()` to handle empty arrays properly (Tom Lane)

• Fix one-byte buffer overrun in `quote_literal_cstr()` (Heikki Linnakangas)
  
The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.

• Prevent multiple calls of `pg_start_backup()` and `pg_stop_backup()` from running concurrently (Michael Paquier)
  
  This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

• Avoid discarding `interval-to-interval` casts that aren't really no-ops (Tom Lane)
  
  In some cases, a cast that should result in zeroing out low-order `interval` fields was mistakenly deemed to be a no-op and discarded. An example is that casting from `INTERVAL MONTH` to `INTERVAL YEAR` failed to clear the months field.

• Ensure that cached plans are invalidated by changes in foreign-table options (Amit Langote, Etsuro Fujita, Ashutosh Bapat)

• Fix `pg_dump` to dump user-defined casts and transforms that use built-in functions (Stephen Frost)

• Fix `pg_restore` with `--create --if-exists` to behave more sanely if an archive contains unrecognized `DROP` commands (Tom Lane)
  
  This doesn't fix any live bug, but it may improve the behavior in future if `pg_restore` is used with an archive generated by a later `pg_dump` version.

• Fix `pg_basebackup`'s rate limiting in the presence of slow I/O (Antonin Houska)
  
  If disk I/O was transiently much slower than the specified rate limit, the calculation overflowed, effectively disabling the rate limit for the rest of the run.

• Fix `pg_basebackup`'s handling of symlinked `pg_stat_tmp` and `pg_replslot` subdirectories (Magnus Hagander, Michael Paquier)

• Fix possible `pg_basebackup` failure on standby server when including WAL files (Amit Kapila, Robert Haas)

• Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)
  
  This avoids failures if the objects are used after a Python garbage collection cycle has occurred.

• Fix PL/Tcl to support triggers on tables that have `.tupno` as a column name (Tom Lane)
  
  This matches the (previously undocumented) behavior of PL/Tcl's `spi_exec` and `spi_execp` commands, namely that a magic `.tupno` column is inserted only if there isn't a real column named that.

• Allow DOS-style line endings in `~/.pgpass` files, even on Unix (Vik Fearing)
  
  This change simplifies use of the same password file across Unix and Windows machines.

• Fix one-byte buffer overrun if `ecpg` is given a file name that ends with a dot (Takayuki Tsunakawa)

• Fix `psql`'s tab completion for `ALTER DEFAULT PRIVILEGES` (Gilles Darold, Stephen Frost)

• In `psql`, treat an empty or all-blank setting of the `PAGER` environment variable as meaning “no pager” (Tom Lane)
  
  Previously, such a setting caused output intended for the pager to vanish entirely.
• Improve contrib/dblink’s reporting of low-level libpq errors, such as out-of-memory (Joe Conway)

• Teach contrib/dblink to ignore irrelevant server options when it uses a contrib/postgres_fdw foreign server as the source of connection options (Corey Huinker)

Previously, if the foreign server object had options that were not also libpq connection options, an error occurred.

• On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)

• Sync our copy of the timezone library with IANA release tzcode2016j (Tom Lane)

This fixes various issues, most notably that timezone data installation failed if the target directory didn’t support hard links.

• Update time zone data files to tzdata release 2016j for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarctica/Casey. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.

E.48. Release 9.4.10

Release date: 2016-10-27

This release contains a variety of fixes from 9.4.9. For information about new features in the 9.4 major release, see Section E.58.

E.48.1. Migration to Version 9.4.10

A dump/restore is not required for those running 9.4.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted free space maps.

Also, if you are upgrading from a version earlier than 9.4.6, see Section E.52.

E.48.2. Changes

• Fix WAL-logging of truncation of relation free space maps and visibility maps (Pavan Deolasee, Heikki Linnakangas)

It was possible for these files to not be correctly restored during crash recovery, or to be written incorrectly on a standby server. Bogus entries in a free space map could lead to attempts to access pages that have been truncated away from the relation itself, typically producing errors like “could not read block XXX: read only 0 of 8192 bytes”. Checksum failures in the visibility map are also possible, if checksumming is enabled.

Procedures for determining whether there is a problem and repairing it if so are discussed at https://wiki.postgresql.org/wiki/Free_Space_Map_Problems.

• Fix incorrect creation of GIN index WAL records on big-endian machines (Tom Lane)

The typical symptom was “unexpected GIN leaf action” errors during WAL replay.

• Fix SELECT FOR UPDATE/SHARE to correctly lock tuples that have been updated by a subsequently-aborted transaction (Álvaro Herrera)
In 9.5 and later, the `SELECT` would sometimes fail to return such tuples at all. A failure has not been proven to occur in earlier releases, but might be possible with concurrent updates.

- Fix EvalPlanQual rechecks involving CTE scans (Tom Lane)

  The recheck would always see the CTE as returning no rows, typically leading to failure to update rows that were recently updated.

- Fix improper repetition of previous results from hashed aggregation in a subquery (Andrew Gierth)

  The test to see if we can reuse a previously-computed hash table of the aggregate state values neglected the possibility of an outer query reference appearing in an aggregate argument expression. A change in the value of such a reference should lead to recalculating the hash table, but did not.

- Fix query-lifespan memory leak in a bulk UPDATE on a table with a PRIMARY KEY or REPLICA IDENTITY index (Tom Lane)

- Fix `EXPLAIN` to emit valid XML when track io timing is on (Markus Winand)

  Previously the XML output-format option produced syntactically invalid tags such as `<I/O-Read-Time>`. That is now rendered as `<I-O-Read-Time>`.

- Suppress printing of zeroes for unmeasured times in `EXPLAIN` (Maksim Milyutin)

  Certain option combinations resulted in printing zero values for times that actually aren't ever measured in that combination. Our general policy in `EXPLAIN` is not to print such fields at all, so do that consistently in all cases.

- Fix timeout length when VACUUM is waiting for exclusive table lock so that it can truncate the table (Simon Riggs)

  The timeout was meant to be 50 milliseconds, but it was actually only 50 microseconds, causing VACUUM to give up on truncation much more easily than intended. Set it to the intended value.

- Fix bugs in merging inherited CHECK constraints while creating or altering a table (Tom Lane, Amit Langote)

  Allow identical CHECK constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a NOT VALID constraint on the child. Likewise, prevent merging of a NO INHERIT child constraint with an inherited constraint.

- Remove artificial restrictions on the values accepted by `numeric_in()` and `numeric_recv()` (Tom Lane)

  We allow numeric values up to the limit of the storage format (more than 1e100000), so it seems fairly pointless that `numeric_in()` rejected scientific-notation exponents above 1000. Likewise, it was silly for `numeric_recv()` to reject more than 1000 digits in an input value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)

- Fix logical WAL decoding to work properly when a subtransaction's WAL output is large enough to spill to disk (Andres Freund)

- Fix buffer overread in logical WAL decoding (Tom Lane)

  Logical decoding of a tuple update record read 23 bytes too many, which was usually harmless but with very bad luck could result in a crash.

- Fix file descriptor leakage when truncating a temporary relation of more than 1GB (Andres Freund)
• Disallow starting a standalone backend with standby_mode turned on (Michael Paquier)

This can’t do anything useful, since there will be no WAL receiver process to fetch more WAL data; and it could result in misbehavior in code that wasn’t designed with this situation in mind.

• Properly initialize replication slot state when recycling a previously-used slot (Michael Paquier)

This failure to reset all of the fields of the slot could prevent VACUUM from removing dead tuples.

• Round shared-memory allocation request to a multiple of the actual huge page size when attempting to use huge pages on Linux (Tom Lane)

This avoids possible failures during munmap() on systems with atypical default huge page sizes. Except in crash-recovery cases, there were no ill effects other than a log message.

• Use a more random value for the dynamic shared memory control segment’s ID (Robert Haas, Tom Lane)

Previously, the same value would be chosen every time, because it was derived from random() but srandom() had not yet been called. While relatively harmless, this was not the intended behavior.

• On Windows, retry creation of the dynamic shared memory control segment after an access-denied error (Kyotaro Horiguchi, Amit Kapila)

Windows sometimes returns ERROR_ACCESS_DENIED rather than ERROR_ALREADY_EXISTS when there is an existing segment. This led to postmaster startup failure due to believing that the former was an unrecoverable error.

• Don’t try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)

This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)

The reported problem involved leaking an error report during PQreset(), but there might be related cases.

• Make ecpg’s --help and --version options work consistently with our other executables (Haribabu Kommi)

• Fix pgbench’s calculation of average latency (Fabien Coelho)

The calculation was incorrect when there were \sleep commands in the script, or when the test duration was specified in number of transactions rather than total time.

• In pg_dump, never dump range constructor functions (Tom Lane)

This oversight led to pg_upgrade failures with extensions containing range types, due to duplicate creation of the constructor functions.

• In pg_xlogdump, retry opening new WAL segments when using --follow option (Magnus Hagander)

This allows for a possible delay in the server’s creation of the next segment.

• Fix pg_xlogdump to cope with a WAL file that begins with a continuation record spanning more than one page (Pavan Deolasee)

• Fix contrib/pg_buffercache to work when shared_buffers exceeds 256GB (KaiGai Kohei)
- Fix contrib/intarray/bench/bench.pl to print the results of the EXPLAIN it does when given the -e option (Daniel Gustafsson)

- Install TAP test infrastructure so that it's available for extension testing (Craig Ringer)

  When PostgreSQL has been configured with --enable-tap-tests, “make install” will now install the Perl support files for TAP testing where PGXS can find them. This allows non-core extensions to use $(prove_check) without extra tests.

- In MSVC builds, include pg_recvlogical in a client-only installation (MauMau)

- Update Windows time zone mapping to recognize some time zone names added in recent Windows versions (Michael Paquier)

- Prevent failure of obsolete dynamic time zone abbreviations (Tom Lane)

  If a dynamic time zone abbreviation does not match any entry in the referenced time zone, treat it as equivalent to the time zone name. This avoids unexpected failures when IANA removes abbreviations from their time zone database, as they did in tzdata release 2016f and seem likely to do again in the future. The consequences were not limited to not recognizing the individual abbreviation; any mismatch caused the pg_timezone_abbrevs view to fail altogether.

- Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

  The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

  In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

**E.49. Release 9.4.9**

**Release date:** 2016-08-11

This release contains a variety of fixes from 9.4.8. For information about new features in the 9.4 major release, see Section E.58.

**E.49.1. Migration to Version 9.4.9**

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.6, see Section E.52.

**E.49.2. Changes**

- Fix possible mis-evaluation of nested CASE-WHEN expressions (Heikki Linnakangas, Michael Paquier, Tom Lane)

  A CASE expression appearing within the test value subexpression of another CASE could become confused about whether its own test value was null or not. Also, inlining of a SQL function implementing the equality operator used by a CASE expression could result in passing the wrong test value to functions called within a CASE expression in the SQL function's body. If the test values
were of different data types, a crash might result; moreover such situations could be abused to allow
disclosure of portions of server memory. (CVE-2016-5423)

- Fix client programs’ handling of special characters in database and role names (Noah Misch, Nathan
  Bossart, Michael Paquier)

Numerous places in vacuumdb and other client programs could become confused by database and
role names containing double quotes or backslashes. Tighten up quoting rules to make that safe.
Also, ensure that when a conninfo string is used as a database name parameter to these programs,
it is correctly treated as such throughout.

Fix handling of paired double quotes in psql’s \connect and \password commands to match
the documentation.

Introduce a new -reuse-previous option in psql’s \connect command to allow explicit
control of whether to re-use connection parameters from a previous connection. (Without this, the
choice is based on whether the database name looks like a conninfo string, as before.) This allows
secure handling of database names containing special characters in pg_dumpall scripts.

pg_dumpall now refuses to deal with database and role names containing carriage returns or new-
lines, as it seems impractical to quote those characters safely on Windows. In future we may reject
such names on the server side, but that step has not been taken yet.

These are considered security fixes because crafted object names containing special characters could
have been used to execute commands with superuser privileges the next time a superuser executes
gp_dumpall or other routine maintenance operations. (CVE-2016-5424)

- Fix corner-case misbehaviors for IS NULL/IS NOT NULL applied to nested composite values
  (Andrew Gierth, Tom Lane)

The SQL standard specifies that IS NULL should return TRUE for a row of all null values (thus
ROW(NULL,NULL) IS NULL yields TRUE), but this is not meant to apply recursively (thus
ROW(NULL, ROW(NULL,NULL)) IS NULL yields FALSE). The core executor got this right,
but certain planner optimizations treated the test as recursive (thus producing TRUE in both cases),
and contrib/postgres_fdw could produce remote queries that misbehaved similarly.

- Make the inet and cidr data types properly reject IPv6 addresses with too many colon-separated
  fields (Tom Lane)

- Prevent crash in close_ps() (the point # lseg operator) for NaN input coordinates (Tom
  Lane)

Make it return NULL instead of crashing.

- Avoid possible crash in pg_get_expr() when inconsistent values are passed to it (Michael
  Paquier, Thomas Munro)

- Fix several one-byte buffer over-reads in to_number() (Peter Eisentraut)

In several cases the to_number() function would read one more character than it should from
the input string. There is a small chance of a crash, if the input happens to be adjacent to the end
of memory.

- Do not run the planner on the query contained in CREATE MATERIALIZED VIEW or CREATE
  TABLE AS when WITH NO DATA is specified (Michael Paquier, Tom Lane)

This avoids some unnecessary failure conditions, for example if a stable function invoked by the
materialized view depends on a table that doesn’t exist yet.

- Avoid unsafe intermediate state during expensive paths through heap_update() (Masahiko
  Sawada, Andres Freund)
Previously, these cases locked the target tuple (by setting its XMAX) but did not WAL-log that action, thus risking data integrity problems if the page were spilled to disk and then a database crash occurred before the tuple update could be completed.

- Fix hint bit update during WAL replay of row locking operations (Andres Freund)

  The only known consequence of this problem is that row locks held by a prepared, but uncommitted, transaction might fail to be enforced after a crash and restart.

- Avoid unnecessary “could not serialize access” errors when acquiring FOR KEY SHARE row locks in serializable mode (Álvaro Herrera)

- Avoid crash in `postgres -C` when the specified variable has a null string value (Michael Paquier)

- Fix possible loss of large subtransactions in logical decoding (Petru-Florin Mihancea)

- Fix failure of logical decoding when a subtransaction contains no actual changes (Marko Tiikkaja, Andrew Gierth)

- Ensure that backends see up-to-date statistics for shared catalogs (Tom Lane)

  The statistics collector failed to update the statistics file for shared catalogs after a request from a regular backend. This problem was partially masked because the autovacuum launcher regularly makes requests that did cause such updates; however, it became obvious with autovacuum disabled.

- Avoid redundant writes of the statistics files when multiple backends request updates close together (Tom Lane, Tomas Vondra)

- Avoid consuming a transaction ID during VACUUM (Alexander Korotkov)

  Some cases in VACUUM unnecessarily caused an XID to be assigned to the current transaction. Normally this is negligible, but if one is up against the XID wraparound limit, consuming more XIDs during anti-wraparound vacuums is a very bad thing.

- Avoid canceling hot-standby queries during VACUUM FREEZE (Simon Riggs, Álvaro Herrera)

  VACUUM FREEZE on an otherwise-idle master server could result in unnecessary cancellations of queries on its standby servers.

- Prevent possible failure when vacuuming multixact IDs in an installation that has been pg_upgrade'd from pre-9.3 (Andrew Gierth, Álvaro Herrera)

  The usual symptom of this bug is errors like “MultiXactId NNN has not been created yet -- apparent wraparound”.

- When a manual ANALYZE specifies a column list, don't reset the table's changes_since_analyze counter (Tom Lane)

  If we're only analyzing some columns, we should not prevent routine auto-analyze from happening for the other columns.

- Fix ANALYZE's overestimation of n_distinct for a unique or nearly-unique column with many null entries (Tom Lane)

  The nulls could get counted as though they were themselves distinct values, leading to serious planner misestimates in some types of queries.

- Prevent autovacuum from starting multiple workers for the same shared catalog (Álvaro Herrera)

  Normally this isn't much of a problem because the vacuum doesn't take long anyway; but in the case of a severely bloated catalog, it could result in all but one worker uselessly waiting instead of doing useful work on other tables.
• Avoid duplicate buffer lock release when abandoning a b-tree index page deletion attempt (Tom Lane)

This mistake prevented VACUUM from completing in some cases involving corrupt b-tree indexes.

• Prevent infinite loop in GiST index build for geometric columns containing NaN component values (Tom Lane)

• Fix contrib/btree_gin to handle the smallest possible bigint value correctly (Peter Eisentraut)

• Teach libpq to correctly decode server version from future servers (Peter Eisentraut)

It's planned to switch to two-part instead of three-part server version numbers for releases after 9.6. Make sure that PQserverVersion() returns the correct value for such cases.

• Fix ecpg's code for unsigned long long array elements (Michael Meskes)

• In pg_dump with both -c and -C options, avoid emitting an unwanted CREATE SCHEMA public command (David Johnston, Tom Lane)

• Improve handling of SIGTERM/control-C in parallel pg_dump and pg_restore (Tom Lane)

Make sure that the worker processes will exit promptly, and also arrange to send query-cancel requests to the connected backends, in case they are doing something long-running such as a CREATE INDEX.

• Fix error reporting in parallel pg_dump and pg_restore (Tom Lane)

Previously, errors reported by pg_dump or pg_restore worker processes might never make it to the user's console, because the messages went through the master process, and there were various deadlock scenarios that would prevent the master process from passing on the messages. Instead, just print everything to stderr. In some cases this will result in duplicate messages (for instance, if all the workers report a server shutdown), but that seems better than no message.

• Ensure that parallel pg_dump or pg_restore on Windows will shut down properly after an error (Kyotaro Horiguchi)

Previously, it would report the error, but then just sit until manually stopped by the user.

• Make pg_dump behave better when built without zlib support (Kyotaro Horiguchi)

It didn't work right for parallel dumps, and emitted some rather pointless warnings in other cases.

• Make pg_basebackup accept -Z 0 as specifying no compression (Fujii Masao)

• Fix makefiles' rule for building AIX shared libraries to be safe for parallel make (Noah Misch)

• Fix TAP tests and MSVC scripts to work when build directory's path name contains spaces (Michael Paquier, Kyotaro Horiguchi)

• Be more predictable about reporting "statement timeout" versus "lock timeout" (Tom Lane)

On heavily loaded machines, the regression tests sometimes failed due to reporting "lock timeout" even though the statement timeout should have occurred first.

• Make regression tests safe for Danish and Welsh locales (Jeff Janes, Tom Lane)

Change some test data that triggered the unusual sorting rules of these locales.

• Update our copy of the timezone code to match IANA's tzcode release 2016c (Tom Lane)

2574
This is needed to cope with anticipated future changes in the time zone data files. It also fixes some corner-case bugs in coping with unusual time zones.

• Update time zone data files to tzdata release 2016f for DST law changes in Kemerovo and Novosibirsk, plus historical corrections for Azerbaijan, Belarus, and Morocco.

E.50. Release 9.4.8

Release date: 2016-05-12

This release contains a variety of fixes from 9.4.7. For information about new features in the 9.4 major release, see Section E.58.

E.50.1. Migration to Version 9.4.8

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.6, see Section E.52.

E.50.2. Changes

• Clear the OpenSSL error queue before OpenSSL calls, rather than assuming it’s clear already; and make sure we leave it clear afterwards (Peter Geoghegan, Dave Vitek, Peter Eisentraut)

This change prevents problems when there are multiple connections using OpenSSL within a single process and not all the code involved follows the same rules for when to clear the error queue. Failures have been reported specifically when a client application uses SSL connections in libpq concurrently with SSL connections using the PHP, Python, or Ruby wrappers for OpenSSL. It’s possible for similar problems to arise within the server as well, if an extension module establishes an outgoing SSL connection.

• Fix “failed to build any n-way joins” planner error with a full join enclosed in the right-hand side of a left join (Tom Lane)

• Fix incorrect handling of equivalence-class tests in multilevel nestloop plans (Tom Lane)

Given a three-or-more-way equivalence class of variables, such as \( X.X = Y.Y = Z.Z \), it was possible for the planner to omit some of the tests needed to enforce that all the variables are actually equal, leading to join rows being output that didn’t satisfy the \( \text{WHERE} \) clauses. For various reasons, erroneous plans were seldom selected in practice, so that this bug has gone undetected for a long time.

• Fix query-lifespan memory leak in GIN index scans (Julien Rouhaud)

• Fix query-lifespan memory leak and potential index corruption hazard in GIN index insertion (Tom Lane)

The memory leak would typically not amount to much in simple queries, but it could be very substantial during a large GIN index build with high \text{maintenance\_work\_mem}.

• Fix possible misbehavior of \( \text{TH}, \text{th}, \text{and Y, YYY} \) format codes in \text{to\_timestamp()} (Tom Lane)

These could advance off the end of the input string, causing subsequent format codes to read garbage.

• Fix dumping of rules and views in which the \text{array} argument of a \text{value operator ANY (array)} construct is a sub-SELECT (Tom Lane)

• Disallow newlines in \text{ALTER SYSTEM} parameter values (Tom Lane)
The configuration-file parser doesn’t support embedded newlines in string literals, so we mustn’t allow them in values to be inserted by ALTER SYSTEM.

- Fix ALTER TABLE ... REPLICA IDENTITY USING INDEX to work properly if an index on OID is selected (David Rowley)

- Fix crash in logical decoding on alignment-picky platforms (Tom Lane, Andres Freund)

  The failure occurred only with a transaction large enough to spill to disk and a primary-key change within that transaction.

- Avoid repeated requests for feedback from receiver while shutting down walsender (Nick Cleaton)

- Make pg_regress use a startup timeout from the PGCTLTIMEOUT environment variable, if that’s set (Tom Lane)

  This is for consistency with a behavior recently added to pg_ctl; it eases automated testing on slow machines.

- Fix pg_upgrade to correctly restore extension membership for operator families containing only one operator class (Tom Lane)

  In such a case, the operator family was restored into the new database, but it was no longer marked as part of the extension. This had no immediate ill effects, but would cause later pg_dump runs to emit output that would cause (harmless) errors on restore.

- Fix pg_upgrade to not fail when new-cluster TOAST rules differ from old (Tom Lane)

  pg_upgrade had special-case code to handle the situation where the new PostgreSQL version thinks that a table should have a TOAST table while the old version did not. That code was broken, so remove it, and instead do nothing in such cases; there seems no reason to believe that we can’t get along fine without a TOAST table if that was okay according to the old version’s rules.

- Reduce the number of SysV semaphores used by a build configured with --disable-spinlocks (Tom Lane)

- Rename internal function strtoi() to strtoint() to avoid conflict with a NetBSD library function (Thomas Munro)

- Fix reporting of errors from bind() and listen() system calls on Windows (Tom Lane)

- Reduce verbosity of compiler output when building with Microsoft Visual Studio (Christian Ullrich)

- Fix putenv() to work properly with Visual Studio 2013 (Michael Paquier)

- Avoid possibly-unsafe use of Windows’ FormatMessage() function (Christian Ullrich)

  Use the FORMAT_MESSAGE_IGNORE_INSERTS flag where appropriate. No live bug is known to exist here, but it seems like a good idea to be careful.

- Update time zone data files to tzdata release 2016d for DST law changes in Russia and Venezuela. There are new zone names Europe/Kirov and Asia/Tomsk to reflect the fact that these regions now have different time zone histories from adjacent regions.

E.51. Release 9.4.7

Release date: 2016-03-31

This release contains a variety of fixes from 9.4.6. For information about new features in the 9.4 major release, see Section E.58.
E.51.1. Migration to Version 9.4.7

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.6, see Section E.52.

E.51.2. Changes

- Fix incorrect handling of NULL index entries in indexed \texttt{ROW()} comparisons (Tom Lane)

  An index search using a row comparison such as \texttt{ROW(a, b) > ROW('x', 'y')} would stop upon reaching a NULL entry in the \texttt{b} column, ignoring the fact that there might be non-NULL \texttt{b} values associated with later values of \texttt{a}.

- Avoid unlikely data-loss scenarios due to renaming files without adequate \texttt{fsync()} calls before and after (Michael Paquier, Tomas Vondra, Andres Freund)

- Fix bug in \texttt{json_to_record()} when a field of its input object contains a sub-object with a field name matching one of the requested output column names (Tom Lane)

- Fix misformatting of negative time zone offsets by \texttt{to_char()}'s \texttt{OF} format code (Thomas Munro, Tom Lane)

- Ignore recovery\_min\_apply\_delay parameter until recovery has reached a consistent state (Michael Paquier)

  Previously, standby servers would delay application of WAL records in response to recovery\_min\_apply\_delay even while replaying the initial portion of WAL needed to make their database state valid. Since the standby is useless until it's reached a consistent database state, this was deemed unhelpful.

- Correctly handle cases where \texttt{pg\_subtrans} is close to XID wraparound during server startup (Jeff Janes)

- Fix assorted bugs in logical decoding (Andres Freund)

  Trouble cases included tuples larger than one page when replica identity is \texttt{FULL}, \texttt{UPDATES} that change a primary key within a transaction large enough to be spooled to disk, incorrect reports of "subxact logged without previous toplevel record", and incorrect reporting of a transaction's commit time.

- Fix planner error with nested security barrier views when the outer view has a \texttt{WHERE} clause containing a correlated subquery (Dean Rasheed)

- Fix corner-case crash due to trying to free \texttt{localeconv()} output strings more than once (Tom Lane)

- Fix parsing of affix files for \texttt{ispell} dictionaries (Tom Lane)

  The code could go wrong if the affix file contained any characters whose byte length changes during case-folding, for example \texttt{I} in Turkish UTF8 locales.

- Avoid use of \texttt{sscanf()} to parse \texttt{ispell} dictionary files (Artur Zakirov)

  This dodges a portability problem on FreeBSD-derived platforms (including macOS).

- Avoid a crash on old Windows versions (before 7SP1/2008R2SP1) with an AVX2-capable CPU and a Postgres build done with Visual Studio 2013 (Christian Ullrich)

  This is a workaround for a bug in Visual Studio 2013's runtime library, which Microsoft have stated they will not fix in that version.
• Fix psql's tab completion logic to handle multibyte characters properly (Kyotaro Horiguchi, Robert Haas)

• Fix psql's tab completion for SECURITY LABEL (Tom Lane)
  Pressing TAB after SECURITY LABEL might cause a crash or offering of inappropriate keywords.

• Make pg_ctl accept a wait timeout from the PGCTLTIMEOUT environment variable, if none is specified on the command line (Noah Misch)
  This eases testing of slower buildfarm members by allowing them to globally specify a longer-than-normal timeout for postmaster startup and shutdown.

• Fix incorrect test for Windows service status in pg_ctl (Manuel Mathar)
  The previous set of minor releases attempted to fix pg_ctl to properly determine whether to send log messages to Window's Event Log, but got the test backwards.

• Fix pgbench to correctly handle the combination of -C and -M prepared options (Tom Lane)

• In pg_upgrade, skip creating a deletion script when the new data directory is inside the old data directory (Bruce Momjian)
  Blind application of the script in such cases would result in loss of the new data directory.

• In PL/Perl, properly translate empty Postgres arrays into empty Perl arrays (Alex Hunsaker)

• Make PL/Python cope with function names that aren't valid Python identifiers (Jim Nasby)

• Fix multiple mistakes in the statistics returned by contrib/pgstattuple's pgstatindex() function (Tom Lane)

• Remove dependency on psed in MSVC builds, since it's no longer provided by core Perl (Michael Paquier, Andrew Dunstan)

• Update time zone data files to tzdata release 2016c for DST law changes in Azerbaijan, Chile, Haiti, Palestine, and Russia (Altai, Astrakhan, Kirov, Sakhalin, Ulyanovsk regions), plus historical corrections for Lithuania, Moldova, and Russia (Kaliningrad, Samara, Volgograd).

E.52. Release 9.4.6

Release date: 2016-02-11

This release contains a variety of fixes from 9.4.5. For information about new features in the 9.4 major release, see Section E.58.

E.52.1. Migration to Version 9.4.6

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading an installation that contains any GIN indexes that use the (non-default) jsonb_path_ops operator class, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.4.4, see Section E.54.

E.52.2. Changes

• Fix inconsistent hash calculations in jsonb_path_ops GIN indexes (Tom Lane)
  When processing jsonb values that contain both scalars and sub-objects at the same nesting level, for example an array containing both scalars and sub-arrays, key hash values could be calculated
differently than they would be for the same key in a different context. This could result in queries not finding entries that they should find. Fixing this means that existing indexes may now be inconsistent with the new hash calculation code. Users should REINDEX jsonb_path_ops GIN indexes after installing this update to make sure that all searches work as expected.

- Fix infinite loops and buffer-overrun problems in regular expressions (Tom Lane)

Very large character ranges in bracket expressions could cause infinite loops in some cases, and memory overwrites in other cases. (CVE-2016-0773)

- Perform an immediate shutdown if the postmaster.pid file is removed (Tom Lane)

The postmaster now checks every minute or so that postmaster.pid is still there and still contains its own PID. If not, it performs an immediate shutdown, as though it had received SIGQUIT. The main motivation for this change is to ensure that failed buildfarm runs will get cleaned up without manual intervention; but it also serves to limit the bad effects if a DBA forcibly removes postmaster.pid and then starts a new postmaster.

- In Serializable transaction isolation mode, serialization anomalies could be missed due to race conditions during insertions (Kevin Grittner, Thomas Munro)

- Fix failure to emit appropriate WAL records when doing ALTER TABLE ... SET TABLESPACE for unlogged relations (Michael Paquier, Andres Freund)

Even though the relation's data is unlogged, the move must be logged or the relation will be inaccessible after a standby is promoted to master.

- Fix possible misinitialization of unlogged relations at the end of crash recovery (Andres Freund, Michael Paquier)

- Ensure walsender slots are fully re-initialized when being re-used (Magnus Hagander)

- Fix ALTER COLUMN TYPE to reconstruct inherited check constraints properly (Tom Lane)

- Fix REASSIGN OWNED to change ownership of composite types properly (Álvaro Herrera)

- Fix REASSIGN OWNED and ALTER OWNER to correctly update granted-permissions lists when changing owners of data types, foreign data wrappers, or foreign servers (Bruce Momjian, Álvaro Herrera)

- Fix REASSIGN OWNED to ignore foreign user mappings, rather than fail (Álvaro Herrera)

- Fix possible crash after doing query rewrite for an updatable view (Stephen Frost)

- Fix planner's handling of LATERAL references (Tom Lane)

This fixes some corner cases that led to “failed to build any N-way joins” or “could not devise a query plan” planner failures.

- Add more defenses against bad planner cost estimates for GIN index scans when the index's internal statistics are very out-of-date (Tom Lane)

- Make planner cope with hypothetical GIN indexes suggested by an index advisor plug-in (Julien Rouhaud)

- Speed up generation of unique table aliases in EXPLAIN and rule dumping, and ensure that generated aliases do not exceed NAMEDATALEN (Tom Lane)

- Fix dumping of whole-row Vars in ROW() and VALUES() lists (Tom Lane)

- Translation of minus-infinity dates and timestamps to json or jsonb incorrectly rendered them as plus-infinity (Tom Lane)
• Fix possible internal overflow in numeric division (Dean Rasheed)

• Fix enforcement of restrictions inside parentheses within regular expression lookahead constraints (Tom Lane)

  Lookahead constraints aren’t allowed to contain backrefs, and parentheses within them are always considered non-capturing, according to the manual. However, the code failed to handle these cases properly inside a parenthesized subexpression, and would give unexpected results.

• Conversion of regular expressions to indexscan bounds could produce incorrect bounds from reg-exps containing lookahead constraints (Tom Lane)

• Fix regular-expression compiler to handle loops of constraint arcs (Tom Lane)

  The code added for CVE-2007-4772 was both incomplete, in that it didn’t handle loops involving more than one state, and incorrect, in that it could cause assertion failures (though there seem to be no bad consequences of that in a non-assert build). Multi-state loops would cause the compiler to run until the query was canceled or it reached the too-many-states error condition.

• Improve memory-usage accounting in regular-expression compiler (Tom Lane)

  This causes the code to emit “regular expression is too complex” errors in some cases that previously used unreasonable amounts of time and memory.

• Improve performance of regular-expression compiler (Tom Lane)

• Make %h and %r escapes in log_line_prefix work for messages emitted due to log_connections (Tom Lane)

  Previously, %h/%r started to work just after a new session had emitted the “connection received” log message; now they work for that message too.

• On Windows, ensure the shared-memory mapping handle gets closed in child processes that don't need it (Tom Lane, Amit Kapila)

  This oversight resulted in failure to recover from crashes whenever logging_collector is turned on.

• Fix possible failure to detect socket EOF in non-blocking mode on Windows (Tom Lane)

  It’s not entirely clear whether this problem can happen in pre-9.5 branches, but if it did, the symptom would be that a walsender process would wait indefinitely rather than noticing a loss of connection.

• Avoid leaking a token handle during SSPI authentication (Christian Ullrich)

• In psql, ensure that libreadline’s idea of the screen size is updated when the terminal window size changes (Merlin Moncure)

  Previously, libreadline did not notice if the window was resized during query output, leading to strange behavior during later input of multiline queries.

• Fix psql’s \det command to interpret its pattern argument the same way as other \d commands with potentially schema-qualified patterns do (Reece Hart)

• Avoid possible crash in psql’s \c command when previous connection was via Unix socket and command specifies a new hostname and same username (Tom Lane)

• In pg_ctl start -w, test child process status directly rather than relying on heuristics (Tom Lane, Michael Paquier)

  Previously, pg_ctl relied on an assumption that the new postmaster would always create postmaster.pid within five seconds. But that can fail on heavily-loaded systems, causing pg_ctl to report incorrectly that the postmaster failed to start.
Except on Windows, this change also means that a `pg_ctl start -w` done immediately after another such command will now reliably fail, whereas previously it would report success if done within two seconds of the first command.

- In `pg_ctl start -w`, don't attempt to use a wildcard listen address to connect to the postmaster (Kondo Yuta)

  On Windows, `pg_ctl` would fail to detect postmaster startup if `listen_addresses` is set to `0.0.0.0` or `::`, because it would try to use that value verbatim as the address to connect to, which doesn't work. Instead assume that `127.0.0.1` or `::1`, respectively, is the right thing to use.

- In `pg_ctl` on Windows, check service status to decide where to send output, rather than checking if standard output is a terminal (Michael Paquier)

- In `pg_dump` and `pg_basebackup`, adopt the GNU convention for handling tar-archive members exceeding 8GB (Tom Lane)

  The POSIX standard for tar file format does not allow archive member files to exceed 8GB, but most modern implementations of tar support an extension that fixes that. Adopt this extension so that `pg_dump` with `-Ft` no longer fails on tables with more than 8GB of data, and so that `pg_basebackup` can handle files larger than 8GB. In addition, fix some portability issues that could cause failures for members between 4GB and 8GB on some platforms. Potentially these problems could cause unrecoverable data loss due to unreadable backup files.

- Fix assorted corner-case bugs in `pg_dump`'s processing of extension member objects (Tom Lane)

- Make `pg_dump` mark a view's triggers as needing to be processed after its rule, to prevent possible failure during parallel `pg_restore` (Tom Lane)

- Ensure that relation option values are properly quoted in `pg_dump` (Kouhei Sutou, Tom Lane)

  A reloption value that isn't a simple identifier or number could lead to dump/reload failures due to syntax errors in CREATE statements issued by `pg_dump`. This is not an issue with any reloption currently supported by core PostgreSQL, but extensions could allow reloptions that cause the problem.

- Avoid repeated password prompts during parallel `pg_dump` (Zeus Kronion)

- Fix `pg_upgrade`'s file-copying code to handle errors properly on Windows (Bruce Momjian)

- Install guards in pgbench against corner-case overflow conditions during evaluation of script-specified division or modulo operators (Fabien Coelho, Michael Paquier)

- Fix failure to localize messages emitted by `pg_recv*` and `pg_receiv*` (Ioseph Kim)

- Avoid dump/reload problems when using both plpython2 and plpython3 (Tom Lane)

  In principle, both versions of PL/Python can be used in the same database, though not in the same session (because the two versions of libpython cannot safely be used concurrently). However, `pg_restore` and `pg_upgrade` both do things that can fall foul of the same-session restriction. Work around that by changing the timing of the check.

- Fix PL/Python regression tests to pass with Python 3.5 (Peter Eisentraut)

- Fix premature clearing of libpq's input buffer when socket EOF is seen (Tom Lane)

  This mistake caused libpq to sometimes not report the backend's final error message before reporting "server closed the connection unexpectedly".

- Prevent certain PL/Java parameters from being set by non-superusers (Noah Misch)
This change mitigates a PL/Java security bug (CVE-2016-0766), which was fixed in PL/Java by marking these parameters as superuser-only. To fix the security hazard for sites that update PostgreSQL more frequently than PL/Java, make the core code aware of them also.

- Improve libpq's handling of out-of-memory situations (Michael Paquier, Amit Kapila, Heikki Linnakangas)
- Fix order of arguments in ecpg-generated typedef statements (Michael Meskes)
- Use %g not %f format in ecpg's PGTYPESnumeric_from_double() (Tom Lane)
- Fix ecpg-supplied header files to not contain comments continued from a preprocessor directive line onto the next line (Michael Meskes)

Such a comment is rejected by ecpg. It's not yet clear whether ecpg itself should be changed.

- Fix hstore_to_json_loose()'s test for whether an hstore value can be converted to a JSON number (Tom Lane)

Previously this function could be fooled by non-alphanumeric trailing characters, leading to emitting syntactically-invalid JSON.

- Ensure that contrib/pgcrypto's crypt() function can be interrupted by query cancel (Andreas Karlsson)
- In contrib/postgres_fdw, fix bugs triggered by use of tableoid in data-modifying commands (Etsuro Fujita, Robert Haas)
- Accept flex versions later than 2.5.x (Tom Lane, Michael Paquier)

Now that flex 2.6.0 has been released, the version checks in our build scripts needed to be adjusted.

- Improve reproducibility of build output by ensuring filenames are given to the linker in a fixed order (Christoph Berg)

This avoids possible bitwise differences in the produced executable files from one build to the next.

- Install our missing script where PGXS builds can find it (Jim Nasby)

This allows sane behavior in a PGXS build done on a machine where build tools such as bison are missing.

- Ensure that dynloader.h is included in the installed header files in MSVC builds (Bruce Momjian, Michael Paquier)
- Add variant regression test expected-output file to match behavior of current libxml2 (Tom Lane)

The fix for libxml2's CVE-2015-7499 causes it not to output error context reports in some cases where it used to do so. This seems to be a bug, but we'll probably have to live with it for some time, so work around it.

- Update time zone data files to tzdata release 2016a for DST law changes in Cayman Islands, Metlakatla, and Trans-Baikal Territory (Zabaykalsky Krai), plus historical corrections for Pakistan.

E.53. Release 9.4.5

Release date: 2015-10-08

This release contains a variety of fixes from 9.4.4. For information about new features in the 9.4 major release, see Section E.58.
E.53.1. Migration to Version 9.4.5

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.4, see Section E.54.

E.53.2. Changes

• Guard against stack overflows in `json` parsing (Oskari Saarenmaa)

  If an application constructs PostgreSQL `json` or `jsonb` values from arbitrary user input, the application's users can reliably crash the PostgreSQL server, causing momentary denial of service. (CVE-2015-5289)

• Fix `contrib/pgcrypto` to detect and report too-short `crypt()` salts (Josh Kupershmidt)

  Certain invalid salt arguments crashed the server or disclosed a few bytes of server memory. We have not ruled out the viability of attacks that arrange for presence of confidential information in the disclosed bytes, but they seem unlikely. (CVE-2015-5288)

• Fix subtransaction cleanup after a portal (cursor) belonging to an outer subtransaction fails (Tom Lane, Michael Paquier)

  A function executed in an outer-subtransaction cursor could cause an assertion failure or crash by referencing a relation created within an inner subtransaction.

• Fix possible deadlock during WAL insertion when `commit_delay` is set (Heikki Linnakangas)

• Ensure all relations referred to by an updatable view are properly locked during an update statement (Dean Rasheed)

• Fix insertion of relations into the relation cache “init file” (Tom Lane)

  An oversight in a patch in the most recent minor releases caused `pg_trigger_tgrelid_tgname_index` to be omitted from the init file. Subsequent sessions detected this, then deemed the init file to be broken and silently ignored it, resulting in a significant degradation in session startup time. In addition to fixing the bug, install some guards so that any similar future mistake will be more obvious.

• Avoid O(N^2) behavior when inserting many tuples into a SPI query result (Neil Conway)

• Improve `LISTEN` startup time when there are many unread notifications (Matt Newell)

• Fix performance problem when a session alters large numbers of foreign key constraints (Jan Wieck, Tom Lane)

  This was seen primarily when restoring `pg_dump` output for databases with many thousands of tables.

• Disable SSL renegotiation by default (Michael Paquier, Andres Freund)

  While use of SSL renegotiation is a good idea in theory, we have seen too many bugs in practice, both in the underlying OpenSSL library and in our usage of it. Renegotiation will be removed entirely in 9.5 and later. In the older branches, just change the default value of `ssl_renegotiation_limit` to zero (disabled).

• Lower the minimum values of the `*_freeze_max_age` parameters (Andres Freund)

  This is mainly to make tests of related behavior less time-consuming, but it may also be of value for installations with limited disk space.
• Limit the maximum value of `wal_buffers` to 2GB to avoid server crashes (Josh Berkus)

• Avoid logging complaints when a parameter that can only be set at server start appears multiple times in `postgresql.conf`, and fix counting of line numbers after an `include_dir` directive (Tom Lane)

• Fix rare internal overflow in multiplication of numeric values (Dean Rasheed)

• Guard against hard-to-reach stack overflows involving record types, range types, `json`, `jsonb`, `tsquery`, `ltxtquery`, and `query_int` (Noah Misch)

• Fix handling of `DOW` and `DOY` in datetime input (Greg Stark)

  These tokens aren’t meant to be used in datetime values, but previously they resulted in opaque internal error messages rather than “invalid input syntax”.

• Add more query-cancel checks to regular expression matching (Tom Lane)

• Add recursion depth protections to regular expression, `SIMILAR TO`, and `LIKE` matching (Tom Lane)

  Suitable search patterns and a low stack depth limit could lead to stack-overrun crashes.

• Fix potential infinite loop in regular expression execution (Tom Lane)

  A search pattern that can apparently match a zero-length string, but actually doesn't match because of a back reference, could lead to an infinite loop.

• In regular expression execution, correctly record match data for capturing parentheses within a quantifier even when the match is zero-length (Tom Lane)

• Fix low-memory failures in regular expression compilation (Andreas Seltenreich)

• Fix low-probability memory leak during regular expression execution (Tom Lane)

• Fix rare low-memory failure in lock cleanup during transaction abort (Tom Lane)

• Fix “unexpected out-of-memory situation during sort” errors when using tuplestores with small `work_mem` settings (Tom Lane)

• Fix very-low-probability stack overrun in `qsort` (Tom Lane)

• Fix “invalid memory alloc request size” failure in hash joins with large `work_mem` settings (Tomas Vondra, Tom Lane)

• Fix assorted planner bugs (Tom Lane)

  These mistakes could lead to incorrect query plans that would give wrong answers, or to assertion failures in assert-enabled builds, or to odd planner errors such as “could not devise a query plan for the given query”, “could not find pathkey item to sort”, “plan should not reference subplan's variable”, or “failed to assign all NestLoopParams to plan nodes”. Thanks are due to Andreas Seltenreich and Piotr Stefaniak for fuzz testing that exposed these problems.

• Improve planner's performance for `UPDATE/DELETE` on large inheritance sets (Tom Lane, Dean Rasheed)

• Ensure standby promotion trigger files are removed at postmaster startup (Michael Paquier, Fujii Masao)

  This prevents unwanted promotion from occurring if these files appear in a database backup that is used to initialize a new standby server.
• During postmaster shutdown, ensure that per-socket lock files are removed and listen sockets are closed before we remove the postmaster.pid file (Tom Lane)

This avoids race-condition failures if an external script attempts to start a new postmaster as soon as pg_ctl stop returns.

• Ensure that the postmaster does not exit until all its child processes are gone, even in an immediate shutdown (Tom Lane)

Like the previous item, this avoids possible race conditions against a subsequently-started postmaster.

• Fix postmaster's handling of a startup-process crash during crash recovery (Tom Lane)

If, during a crash recovery cycle, the startup process crashes without having restored database consistency, we'd try to launch a new startup process, which typically would just crash again, leading to an infinite loop.

• Make emergency autovacuuming for multixact wraparound more robust (Andres Freund)

• Do not print a WARNING when an autovacuum worker is already gone when we attempt to signal it, and reduce log verbosity for such signals (Tom Lane)

• Prevent autovacuum launcher from sleeping unduly long if the server clock is moved backwards a large amount (Álvaro Herrera)

• Ensure that cleanup of a GIN index's pending-insertions list is interruptable by cancel requests (Jeff Janes)

• Allow all-zeroes pages in GIN indexes to be reused (Heikki Linnakangas)

Such a page might be left behind after a crash.

• Fix handling of all-zeroes pages in SP-GiST indexes (Heikki Linnakangas)

VACUUM attempted to recycle such pages, but did so in a way that wasn't crash-safe.

• Fix off-by-one error that led to otherwise-harmless warnings about "apparent wraparound" in subtrans/multixact truncation (Thomas Munro)

• Fix misreporting of CONTINUE and MOVE statement types in PL/pgSQL's error context messages (Pavel Stehule, Tom Lane)

• Fix PL/Perl to handle non-ASCII error message texts correctly (Alex Hunsaker)

• Fix PL/Python crash when returning the string representation of a result (Tom Lane)

• Fix some places in PL/Tcl that neglected to check for failure of malloc() calls (Michael Paquier, Álvaro Herrera)

• In contrib/isn, fix output of ISBN-13 numbers that begin with 979 (Fabien Coelho)

EANs beginning with 979 (but not 9790) are considered ISBNs, but they must be printed in the new 13-digit format, not the 10-digit format.

• Improve contrib/pg_stat_statements' handling of query-text garbage collection (Peter Geoghegan)

The external file containing query texts could bloat to very large sizes; once it got past 1GB attempts to trim it would fail, soon leading to situations where the file could not be read at all.

• Improve contrib/postgres_fdw's handling of collation-related decisions (Tom Lane)
The main user-visible effect is expected to be that comparisons involving \texttt{varchar} columns will be sent to the remote server for execution in more cases than before.

- Improve libpq's handling of out-of-memory conditions (Michael Paquier, Heikki Linnakangas)
- Fix memory leaks and missing out-of-memory checks in ecpg (Michael Paquier)
- Fix psql's code for locale-aware formatting of numeric output (Tom Lane)

The formatting code invoked by \texttt{\textbackslash pset numericlocale on} did the wrong thing for some uncommon cases such as numbers with an exponent but no decimal point. It could also mangle already-localized output from the \texttt{money} data type.

- Prevent crash in psql's \texttt{\textbackslash e} command when there is no current connection (Noah Misch)
- Make \texttt{pg\_dump} handle inherited \texttt{NOT VALID} check constraints correctly (Tom Lane)
- Fix selection of default zlib compression level in \texttt{pg\_dump}'s directory output format (Andrew Dunstan)
- Ensure that temporary files created during a \texttt{pg\_dump} run with tar-format output are not world-readable (Michael Paquier)
- Fix \texttt{pg\_dump} and \texttt{pg\_upgrade} to support cases where the \texttt{postgres} or \texttt{template1} database is in a non-default tablespace (Marti Raudsepp, Bruce Momjian)
- Fix \texttt{pg\_dump} to handle object privileges sanely when dumping from a server too old to have a particular privilege type (Tom Lane)

When dumping data types from pre-9.2 servers, and when dumping functions or procedural languages from pre-7.3 servers, \texttt{pg\_dump} would produce \texttt{GRANT/REVOKE} commands that revoked the owner's grantable privileges and instead granted all privileges to \texttt{PUBLIC}. Since the privileges involved are just \texttt{USAGE} and \texttt{EXECUTE}, this isn't a security problem, but it's certainly a surprising representation of the older systems' behavior. Fix it to leave the default privilege state alone in these cases.

- Fix \texttt{pg\_dump} to dump shell types (Tom Lane)

Shell types (that is, not-yet-fully-defined types) aren't useful for much, but nonetheless \texttt{pg\_dump} should dump them.

- Fix assorted minor memory leaks in \texttt{pg\_dump} and other client-side programs (Michael Paquier)
- Fix pgbench's progress-report behavior when a query, or pgbench itself, gets stuck (Fabien Coelho)
- Fix spinlock assembly code for Alpha hardware (Tom Lane)
- Fix spinlock assembly code for PPC hardware to be compatible with AIX's native assembler (Tom Lane)

Building with gcc didn't work if gcc had been configured to use the native assembler, which is becoming more common.

- On AIX, test the \texttt{-qlonglong} compiler option rather than just assuming it's safe to use (Noah Misch)
- On AIX, use \texttt{-Wl,-brtllib} link option to allow symbols to be resolved at runtime (Noah Misch)

Perl relies on this ability in 5.8.0 and later.
• Avoid use of inline functions when compiling with 32-bit xlc, due to compiler bugs (Noah Misch)

• Use _librt_ for _sched_yield_() when necessary, which it is on some Solaris versions (Oskari Saarenmaa)

• Translate encoding _UHC_ as Windows code page 949 (Noah Misch)

  This fixes presentation of non-ASCII log messages from processes that are not attached to any particular database, such as the postmaster.

• On Windows, avoid failure when doing encoding conversion to UTF16 outside a transaction, such as for log messages (Noah Misch)

• Fix postmaster startup failure due to not copying _setlocale_()'s return value (Noah Misch)

  This has been reported on Windows systems with the ANSI code page set to CP936 (“Chinese (Simplified, PRC”), and may occur with other multibyte code pages.

• Fix Windows _install.bat_ script to handle target directory names that contain spaces (Heikki Linnakangas)

• Make the numeric form of the PostgreSQL version number (e.g., _90405_) readily available to extension Makefiles, as a variable named _VERSION_NUM_ (Michael Paquier)

• Update time zone data files to _tzdata_ release 2015g for DST law changes in Cayman Islands, Fiji, Moldova, Morocco, Norfolk Island, North Korea, Turkey, and Uruguay. There is a new zone name _America/Fort_Nelson_ for the Canadian Northern Rockies.

E.54. Release 9.4.4

**Release date:** 2015-06-12

This release contains a small number of fixes from 9.4.3. For information about new features in the 9.4 major release, see Section E.58.

E.54.1. Migration to Version 9.4.4

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading an installation that was previously upgraded using a _pg_upgrade_ version between 9.3.0 and 9.3.4 inclusive, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.4.2, see Section E.56.

E.54.2. Changes

• Fix possible failure to recover from an inconsistent database state (Robert Haas)

  Recent PostgreSQL releases introduced mechanisms to protect against multixact wraparound, but some of that code did not account for the possibility that it would need to run during crash recovery, when the database may not be in a consistent state. This could result in failure to restart after a crash, or failure to start up a secondary server. The lingering effects of a previously-fixed bug in _pg_upgrade_ could also cause such a failure, in installations that had used _pg_upgrade_ versions between 9.3.0 and 9.3.4.

  The _pg_upgrade_ bug in question was that it would set _oldestMultiXid_ to 1 in _pg_control_ even if the true value should be higher. With the fixes introduced in this release, such a situation will result in immediate emergency autovacuuming until a correct _oldestMultiXid_ value can
be determined. If that would pose a hardship, users can avoid it by doing manual vacuuming before upgrading to this release. In detail:

1. Check whether pg_controldata reports “Latest checkpoint's oldestMultiXid” to be 1. If not, there's nothing to do.

2. Look in PGDATA/pg_multixact/offsets to see if there's a file named 0000. If there is, there's nothing to do.

3. Otherwise, for each table that has pg_class.relminmxid equal to 1, VACUUM that table with both vacuum_multixact_freeze_min_age and vacuum_multixact_freeze_table_age set to zero. (You can use the vacuum cost delay parameters described in Section 19.4.4 to reduce the performance consequences for concurrent sessions.)

• Fix rare failure to invalidate relation cache init file (Tom Lane)

With just the wrong timing of concurrent activity, a VACUUM FULL on a system catalog might fail to update the "init file" that's used to avoid cache-loading work for new sessions. This would result in later sessions being unable to access that catalog at all. This is a very ancient bug, but it's so hard to trigger that no reproducible case had been seen until recently.

• Avoid deadlock between incoming sessions and CREATE/DROP DATABASE (Tom Lane)

A new session starting in a database that is the target of a DROP DATABASE command, or is the template for a CREATE DATABASE command, could cause the command to wait for five seconds and then fail, even if the new session would have exited before that.

• Improve planner's cost estimates for semi-joins and anti-joins with inner indexscans (Tom Lane, Tomas Vondra)

This type of plan is quite cheap when all the join clauses are used as index scan conditions, even if the inner scan would nominally fetch many rows, because the executor will stop after obtaining one row. The planner only partially accounted for that effect, and would therefore overestimate the cost, leading it to possibly choose some other much less efficient plan type.

E.55. Release 9.4.3

Release date: 2015-06-04

This release contains a small number of fixes from 9.4.2. For information about new features in the 9.4 major release, see Section E.58.

E.55.1. Migration to Version 9.4.3

A dump/restore is not required for those running 9.4.X.

However, if you are upgrading from a version earlier than 9.4.2, see Section E.56.

E.55.2. Changes

• Avoid failures while fsync'ing data directory during crash restart (Abhijit Menon-Sen, Tom Lane)

In the previous minor releases we added a patch to fsync everything in the data directory after a crash. Unfortunately its response to any error condition was to fail, thereby preventing the server from starting up, even when the problem was quite harmless. An example is that an unwritable file in the data directory would prevent restart on some platforms; but it is common to make SSL certificate files unwritable by the server. Revise this behavior so that permissions failures are ignored altogether, and other types of failures are logged but do not prevent continuing.
Also apply the same rules in `initdb --sync-only`. This case is less critical but it should act similarly.

- Fix `pg_get_functiondef()` to show functions’ LEAKPROOF property, if set (Jeevan Chalke)
- Fix `pushJsonbValue()` to unpack `jbvBinary` objects (Andrew Dunstan)
  This change does not affect any behavior in the core code as of 9.4, but it avoids a corner case for possible third-party callers.
- Remove configure's check prohibiting linking to a threaded libpython on OpenBSD (Tom Lane)
  The failure this restriction was meant to prevent seems to not be a problem anymore on current OpenBSD versions.

**E.56. Release 9.4.2**

**Release date:** 2015-05-22

This release contains a variety of fixes from 9.4.1. For information about new features in the 9.4 major release, see Section E.58.

**E.56.1. Migration to Version 9.4.2**

A dump/restore is not required for those running 9.4.X.

However, if you use `contrib/citext's regexp_matches()` functions, see the changelog entry below about that.

Also, if you are upgrading from a version earlier than 9.4.1, see Section E.57.

**E.56.2. Changes**

- Avoid possible crash when client disconnects just before the authentication timeout expires (Benkocs Norbert Attila)

  If the timeout interrupt fired partway through the session shutdown sequence, SSL-related state would be freed twice, typically causing a crash and hence denial of service to other sessions. Experimentation shows that an unauthenticated remote attacker could trigger the bug somewhat consistently, hence treat as security issue. (CVE-2015-3165)

- Improve detection of system-call failures (Noah Misch)

  Our replacement implementation of `snprintf()` failed to check for errors reported by the underlying system library calls; the main case that might be missed is out-of-memory situations. In the worst case this might lead to information exposure, due to our code assuming that a buffer had been overwritten when it hadn’t been. Also, there were a few places in which security-relevant calls of other system library functions did not check for failure.

  It remains possible that some calls of the `*printf()` family of functions are vulnerable to information disclosure if an out-of-memory error occurs at just the wrong time. We judge the risk to not be large, but will continue analysis in this area. (CVE-2015-3166)

- In `contrib/pgcrypto`, uniformly report decryption failures as “Wrong key or corrupt data” (Noah Misch)

  Previously, some cases of decryption with an incorrect key could report other error message texts. It has been shown that such variance in error reports can aid attackers in recovering keys from other
systems. While it's unknown whether pgcrypto's specific behaviors are likewise exploitable, it seems better to avoid the risk by using a one-size-fits-all message. (CVE-2015-3167)

- Protect against wraparound of multixact member IDs (Álvaro Herrera, Robert Haas, Thomas Munro)

Under certain usage patterns, the existing defenses against this might be insufficient, allowing pg_multixact/members files to be removed too early, resulting in data loss. The fix for this includes modifying the server to fail transactions that would result in overwriting old multixact member ID data, and improving autovacuum to ensure it will act proactively to prevent multixact member ID wraparound, as it does for transaction ID wraparound.

- Fix incorrect declaration of contrib/citext's regexp_matches() functions (Tom Lane)

These functions should return setof text[], like the core functions they are wrappers for; but they were incorrectly declared as returning just text[]. This mistake had two results: first, if there was no match you got a scalar null result, whereas what you should get is an empty set (zero rows). Second, the g flag was effectively ignored, since you would get only one result array even if there were multiple matches.

While the latter behavior is clearly a bug, there might be applications depending on the former behavior; therefore the function declarations will not be changed by default until PostgreSQL 9.5. In pre-9.5 branches, the old behavior exists in version 1.0 of the citext extension, while we have provided corrected declarations in version 1.1 (which is not installed by default). To adopt the fix in pre-9.5 branches, execute ALTER EXTENSION citext UPDATE TO '1.1' in each database in which citext is installed. (You can also “update” back to 1.0 if you need to undo that.) Be aware that either update direction will require dropping and recreating any views or rules that use citext's regexp_matches() functions.

- Render infinite dates and timestamps as infinity when converting to json, rather than throwing an error (Andrew Dunstan)

- Fix json/jsonb's populate_record() and to_record() functions to handle empty input properly (Andrew Dunstan)

- Fix incorrect checking of deferred exclusion constraints after a HOT update (Tom Lane)

If a new row that potentially violates a deferred exclusion constraint is HOT-updated (that is, no indexed columns change and the row can be stored back onto the same table page) later in the same transaction, the exclusion constraint would be reported as violated when the check finally occurred, even if the row(s) the new row originally conflicted with had been deleted.

- Fix behavior when changing foreign key constraint deferrability status with ALTER TABLE ... ALTER CONSTRAINT (Tom Lane)

Operations later in the same session or concurrent sessions might not honor the status change promptly.

- Fix planning of star-schema-style queries (Tom Lane)

Sometimes, efficient scanning of a large table requires that index parameters be provided from more than one other table (commonly, dimension tables whose keys are needed to index a large fact table). The planner should be able to find such plans, but an overly restrictive search heuristic prevented it.

- Prevent improper reordering of antijoins (NOT EXISTS joins) versus other outer joins (Tom Lane)

This oversight in the planner has been observed to cause “could not find RelOptInfo for given relids” errors, but it seems possible that sometimes an incorrect query plan might get past that consistency check and result in silently-wrong query output.

- Fix incorrect matching of subexpressions in outer-join plan nodes (Tom Lane)
Previously, if textually identical non-strict subexpressions were used both above and below an outer join, the planner might try to re-use the value computed below the join, which would be incorrect because the executor would force the value to NULL in case of an unmatched outer row.

- Fix GEQO planner to cope with failure of its join order heuristic (Tom Lane)

This oversight has been seen to lead to “failed to join all relations together” errors in queries involving LATERAL, and that might happen in other cases as well.

- Ensure that row locking occurs properly when the target of an UPDATE or DELETE is a security-barrier view (Stephen Frost)

- Use a file opened for read/write when syncing replication slot data during database startup (Andres Freund)

On some platforms, the previous coding could result in errors like “could not fsync file "pg_replslot/...": Bad file descriptor”.

- Fix possible deadlock at startup when max_prepared_transactions is too small (Heikki Linnakangas)

- Don’t archive useless preallocated WAL files after a timeline switch (Heikki Linnakangas)

- Recursively f sync() the data directory after a crash (Abhijit Menon-Sen, Robert Haas)

This ensures consistency if another crash occurs shortly later. (The second crash would have to be a system-level crash, not just a database crash, for there to be a problem.)

- Fix autovacuum launcher's possible failure to shut down, if an error occurs after it receives SIGTERM (Álvaro Herrera)

- Fix failure to handle invalidation messages for system catalogs early in session startup (Tom Lane)

This oversight could result in failures in sessions that start concurrently with a VACUUM FULL on a system catalog.

- Fix crash in BackendIdGetTransactionIds() when trying to get status for a backend process that just exited (Tom Lane)

- Cope with unexpected signals in LockBufferForCleanup() (Andres Freund)

This oversight could result in spurious errors about “multiple backends attempting to wait for pin-count 1”.

- Fix crash when doing COPY IN to a table with check constraints that contain whole-row references (Tom Lane)

The known failure case only crashes in 9.4 and up, but there is very similar code in 9.3 and 9.2, so back-patch those branches as well.

- Avoid waiting for WAL flush or synchronous replication during commit of a transaction that was read-only so far as the user is concerned (Andres Freund)

Previously, a delay could occur at commit in transactions that had written WAL due to HOT page pruning, leading to undesirable effects such as sessions getting stuck at startup if all synchronous replicas are down. Sessions have also been observed to get stuck in catchup interrupt processing when using synchronous replication; this will fix that problem as well.

- Avoid busy-waiting with short recovery_min_apply_delay values (Andres Freund)

- Fix crash when manipulating hash indexes on temporary tables (Heikki Linnakangas)
• Fix possible failure during hash index bucket split, if other processes are modifying the index concurrently (Tom Lane)

• Fix memory leaks in GIN index vacuum (Heikki Linnakangas)

• Check for interrupts while analyzing index expressions (Jeff Janes)

  ANALYZE executes index expressions many times; if there are slow functions in such an expression, it’s desirable to be able to cancel the ANALYZE before that loop finishes.

• Ensure tableoid of a foreign table is reported correctly when a READ COMMITTED recheck occurs after locking rows in SELECT FOR UPDATE, UPDATE, or DELETE (Etsuro Fujita)

• Add the name of the target server to object description strings for foreign-server user mappings (Álvaro Herrera)

• Include the schema name in object identity strings for conversions (Álvaro Herrera)

• Recommend setting include_realm to 1 when using Kerberos/GSSAPI/SSPI authentication (Stephen Frost)

  Without this, identically-named users from different realms cannot be distinguished. For the moment this is only a documentation change, but it will become the default setting in PostgreSQL 9.5.

• Remove code for matching IPv4 pg_hba.conf entries to IPv4-in-IPv6 addresses (Tom Lane)

  This hack was added in 2003 in response to a report that some Linux kernels of the time would report IPv4 connections as having IPv4-in-IPv6 addresses. However, the logic was accidentally broken in 9.0. The lack of any field complaints since then shows that it’s not needed anymore. Now we have reports that the broken code causes crashes on some systems, so let’s just remove it rather than fix it. (Had we chosen to fix it, that would make for a subtle and potentially security-sensitive change in the effective meaning of IPv4 pg_hba.conf entries, which does not seem like a good thing to do in minor releases.)

• Fix status reporting for terminated background workers that were never actually started (Robert Haas)

• After a database crash, don’t restart background workers that are marked BGW_NEVER_RESTART (Amit Khandekar)

• Report WAL flush, not insert, position in IDENTIFY_SYSTEM replication command (Heikki Linnakangas)

  This avoids a possible startup failure in pg_receivexlog.

• While shutting down service on Windows, periodically send status updates to the Service Control Manager to prevent it from killing the service too soon; and ensure that pg_ctl will wait for shutdown (Krystian Bigaj)

• Reduce risk of network deadlock when using libpq’s non-blocking mode (Heikki Linnakangas)

  When sending large volumes of data, it’s important to drain the input buffer every so often, in case the server has sent enough response data to cause it to block on output. (A typical scenario is that the server is sending a stream of NOTICE messages during COPY FROM STDIN.) This worked properly in the normal blocking mode, but not so much in non-blocking mode. We’ve modified libpq to opportunistically drain input when it can, but a full defense against this problem requires application cooperation: the application should watch for socket read-ready as well as write-ready conditions, and be sure to call PQconsumeInput() upon read-ready.

• In libpq, fix misparsing of empty values in URI connection strings (Thomas Fanghaenel)
• Fix array handling in ecpg (Michael Meskes)

• Fix psql to sanely handle URIs and conninfo strings as the first parameter to \connect (David Fetter, Andrew Dunstan, Álvaro Herrera)

  This syntax has been accepted (but undocumented) for a long time, but previously some parameters might be taken from the old connection instead of the given string, which was agreed to be undesirable.

• Suppress incorrect complaints from psql on some platforms that it failed to write ~/.psql_history at exit (Tom Lane)

  This misbehavior was caused by a workaround for a bug in very old (pre-2006) versions of libedit. We fixed it by removing the workaround, which will cause a similar failure to appear for anyone still using such versions of libedit. Recommendation: upgrade that library, or use libreadline.

• Fix pg_dump's rule for deciding which casts are system-provided casts that should not be dumped (Tom Lane)

• In pg_dump, fix failure to honor -Z compression level option together with -Fd (Michael Paquier)

• Make pg_dump consider foreign key relationships between extension configuration tables while choosing dump order (Gilles Darold, Michael Paquier, Stephen Frost)

  This oversight could result in producing dumps that fail to reload because foreign key constraints are transiently violated.

• Avoid possible pg_dump failure when concurrent sessions are creating and dropping temporary functions (Tom Lane)

• Fix dumping of views that are just VALUES(...) but have column aliases (Tom Lane)

• Ensure that a view's replication identity is correctly set to nothing during dump/restore (Marko Tiikkaja)

  Previously, if the view was involved in a circular dependency, it might wind up with an incorrect replication identity property.

• In pg_upgrade, force timeline 1 in the new cluster (Bruce Momjian)

  This change prevents upgrade failures caused by bogus complaints about missing WAL history files.

• In pg_upgrade, check for improperly non-connectable databases before proceeding (Bruce Momjian)

• In pg_upgrade, quote directory paths properly in the generated delete_old_cluster script (Bruce Momjian)

• In pg_upgrade, preserve database-level freezing info properly (Bruce Momjian)

  This oversight could cause missing-clog-file errors for tables within the postgres and template1 databases.

• Run pg_upgrade and pg_resetxlog with restricted privileges on Windows, so that they don't fail when run by an administrator (Muhammad Asif Naeeem)

• Improve handling of readdir() failures when scanning directories in initdb and pg_basebackup (Marco Nenciarini)

• Fix slow sorting algorithm in contrib/intarray (Tom Lane)

• Fix compile failure on Sparc V8 machines (Rob Rowan)
• Silence some build warnings on macOS (Tom Lane)

• Update time zone data files to tzdata release 2015d for DST law changes in Egypt, Mongolia, and Palestine, plus historical changes in Canada and Chile. Also adopt revised zone abbreviations for the America/Adak zone (HST/HDT not HAST/HADT).

E.57. Release 9.4.1

Release date: 2015-02-05

This release contains a variety of fixes from 9.4.0. For information about new features in the 9.4 major release, see Section E.58.

E.57.1. Migration to Version 9.4.1

A dump/restore is not required for those running 9.4.X. However, if you are a Windows user and are using the “Norwegian (Bokmål)” locale, manual action is needed after the upgrade to replace any “Norwegian (Bokmål)_Norway” or “norwegian-bokmal” locale names stored in PostgreSQL system catalogs with the plain-ASCII alias “Norwegian_Norway”. For details see http://wiki.postgresql.org/wiki/Changes_To_Norwegian_Locale

E.57.2. Changes

• Fix buffer overruns in \texttt{to\_char()} (Bruce Momjian)

  When \texttt{to\_char()} processes a numeric formatting template calling for a large number of digits, PostgreSQL would read past the end of a buffer. When processing a crafted timestamp formatting template, PostgreSQL would write past the end of a buffer. Either case could crash the server. We have not ruled out the possibility of attacks that lead to privilege escalation, though they seem unlikely. (CVE-2015-0241)

• Fix buffer overrun in replacement *\texttt{printf()} functions (Tom Lane)

  PostgreSQL includes a replacement implementation of \texttt{printf} and related functions. This code will overrun a stack buffer when formatting a floating point number (conversion specifiers \texttt{e}, \texttt{E}, \texttt{f}, \texttt{F}, \texttt{g} or \texttt{G}) with requested precision greater than about 500. This will crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. A database user can trigger such a buffer overrun through the \texttt{to\_char()} SQL function. While that is the only affected core PostgreSQL functionality, extension modules that use printf-family functions may be at risk as well. This issue primarily affects PostgreSQL on Windows. PostgreSQL uses the system implementation of these functions where adequate, which it is on other modern platforms. (CVE-2015-0242)

• Fix buffer overruns in \texttt{contrib/pgcrypto} (Marko Tiikkaja, Noah Misch)

  Errors in memory size tracking within the \texttt{pgcrypto} module permitted stack buffer overruns and improper dependence on the contents of uninitialized memory. The buffer overrun cases can crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. (CVE-2015-0243)

• Fix possible loss of frontend/backend protocol synchronization after an error (Heikki Linnakangas)

  If any error occurred while the server was in the middle of reading a protocol message from the client, it could lose synchronization and incorrectly try to interpret part of the message's data as a new protocol message. An attacker able to submit crafted binary data within a command parameter might succeed in injecting his own SQL commands this way. Statement timeout and query cancellation are the most likely sources of errors triggering this scenario. Particularly vulnerable are applications
that use a timeout and also submit arbitrary user-crafted data as binary query parameters. Disabling statement timeout will reduce, but not eliminate, the risk of exploit. Our thanks to Emil Lenngren for reporting this issue. (CVE-2015-0244)

• Fix information leak via constraint-violation error messages (Stephen Frost)

Some server error messages show the values of columns that violate a constraint, such as a unique constraint. If the user does not have SELECT privilege on all columns of the table, this could mean exposing values that the user should not be able to see. Adjust the code so that values are displayed only when they came from the SQL command or could be selected by the user. (CVE-2014-8161)

• Lock down regression testing's temporary installations on Windows (Noah Misch)

Use SSPI authentication to allow connections only from the OS user who launched the test suite. This closes on Windows the same vulnerability previously closed on other platforms, namely that other users might be able to connect to the test postmaster. (CVE-2014-0067)

• Cope with the Windows locale named “Norwegian (Bokmål)” (Heikki Linnakangas)

Non-ASCII locale names are problematic since it’s not clear what encoding they should be represented in. Map the troublesome locale name to a plain-ASCII alias, “Norwegian_Norway”.

9.4.0 mapped the troublesome name to “norwegian-bokmal”, but that turns out not to work on all Windows configurations. “Norwegian_Norway” is now recommended instead.

• Fix use-of-already-freed-memory problem in EvalPlanQual processing (Tom Lane)

In READ COMMITTED mode, queries that lock or update recently-updated rows could crash as a result of this bug.

• Avoid possible deadlock while trying to acquire tuple locks in EvalPlanQual processing (Álvaro Herrera, Mark Kirkwood)

• Fix failure to wait when a transaction tries to acquire a FOR NO KEY EXCLUSIVE tuple lock, while multiple other transactions currently hold FOR SHARE locks (Álvaro Herrera)

• Improve performance of EXPLAIN with large range tables (Tom Lane)

• Fix jsonb Unicode escape processing, and in consequence disallow \u0000 (Tom Lane)

Previously, the JSON Unicode escape \u0000 was accepted and was stored as those six characters; but that is indistinguishable from what is stored for the input \u0000, resulting in ambiguity. Moreover, in cases where de-escaped textual output is expected, such as the -> operator, the sequence was printed as \u0000, which does not meet the expectation that JSON escaping would be removed. (Consistent behavior would require emitting a zero byte, but PostgreSQL does not support zero bytes embedded in text strings.) 9.4.0 included an ill-advised attempt to improve this situation by adjusting JSON output conversion rules; but of course that could not fix the fundamental ambiguity, and it turned out to break other usages of Unicode escape sequences. Revert that, and to avoid the core problem, reject \u0000 in jsonb input.

If a jsonb column contains a \u0000 value stored with 9.4.0, it will henceforth read out as though it were \u0000, which is the other valid interpretation of the data stored by 9.4.0 for this case.

The json type did not have the storage-ambiguity problem, but it did have the problem of inconsistent de-escaped textual output. Therefore \u0000 will now also be rejected in json values when conversion to de-escaped form is required. This change does not break the ability to store \u0000 in json columns so long as no processing is done on the values. This is exactly parallel to the cases in which non-ASCII Unicode escapes are allowed when the database encoding is not UTF8.

• Fix namespace handling in xpath() (Ali Akbar)
Previously, the xml value resulting from an xpath() call would not have namespace declarations if the namespace declarations were attached to an ancestor element in the input xml value, rather than to the specific element being returned. Propagate the ancestral declaration so that the result is correct when considered in isolation.

- Fix assorted oversights in range-operator selectivity estimation (Emre Hasegeli)

This patch fixes corner-case “unexpected operator NNNN” planner errors, and improves the selectivity estimates for some other cases.

- Revert unintended reduction in maximum size of a GIN index item (Heikki Linnakangas)

9.4.0 could fail with “index row size exceeds maximum” errors for data that previous versions would accept.

- Fix query-duration memory leak during repeated GIN index rescans (Heikki Linnakangas)

- Fix possible crash when using nonzero gin_fuzzy_search_limit (Heikki Linnakangas)

- Assorted fixes for logical decoding (Andres Freund)

- Fix incorrect replay of WAL parameter change records that report changes in the wal_log_hints setting (Petr Jelinek)

- Change “pgstat wait timeout” warning message to be LOG level, and rephrase it to be more understandable (Tom Lane)

  This message was originally thought to be essentially a can’t-happen case, but it occurs often enough on our slower buildfarm members to be a nuisance. Reduce it to LOG level, and expend a bit more effort on the wording: it now reads “using stale statistics instead of current ones because stats collector is not responding”.

- Warn if macOS's setlocale() starts an unwanted extra thread inside the postmaster (Noah Misch)

- Fix libpq's behavior when /etc/passwd isn't readable (Tom Lane)

  While doing PQsetdbLogin(), libpq attempts to ascertain the user's operating system name, which on most Unix platforms involves reading /etc/passwd. As of 9.4, failure to do that was treated as a hard error. Restore the previous behavior, which was to fail only if the application does not provide a database role name to connect as. This supports operation in chroot environments that lack an /etc/passwd file.

- Improve consistency of parsing of psql's special variables (Tom Lane)

  Allow variant spellings of on and off (such as 1/0) for ECHO_HIDDEN and ON_ERROR_ROLLBACK. Report a warning for unrecognized values for COMP_KEYWORD_CASE, ECHO, ECHO_HIDDEN, HISTCONTROL, ON_ERROR_ROLLBACK, and VERBOSITY. Recognize all values for all these variables case-insensitively; previously there was a mishmash of case-sensitive and case-insensitive behaviors.

- Fix pg_dump to handle comments on event triggers without failing (Tom Lane)

- Allow parallel pg_dump to use --serializable-deferrable (Kevin Grittner)

- Prevent WAL files created by pg_basebackup -x/-X from being archived again when the standby is promoted (Andres Freund)

- Handle unexpected query results, especially NULLs, safely in contrib/tablefunc's connectby() (Michael Paquier)
connectby() previously crashed if it encountered a NULL key value. It now prints that row but doesn't recurse further.

• Numerous cleanups of warnings from Coverity static code analyzer (Andres Freund, Tatsuo Ishii, Marko Kreen, Tom Lane, Michael Paquier)

These changes are mostly cosmetic but in some cases fix corner-case bugs, for example a crash rather than a proper error report after an out-of-memory failure. None are believed to represent security issues.

• Allow CFLAGS from configure's environment to override automatically-supplied CFLAGS (Tom Lane)

Previously, configure would add any switches that it chose of its own accord to the end of the user-specified CFLAGS string. Since most compilers process switches left-to-right, this meant that configure's choices would override the user-specified flags in case of conflicts. That should work the other way around, so adjust the logic to put the user's string at the end not the beginning.

• Make pg_regress remove any temporary installation it created upon successful exit (Tom Lane)

This results in a very substantial reduction in disk space usage during make check-world, since that sequence involves creation of numerous temporary installations.

• Add CST (China Standard Time) to our lists of timezone abbreviations (Tom Lane)

• Update time zone data files to tzdata release 2015a for DST law changes in Chile and Mexico, plus historical changes in Iceland.

E.58. Release 9.4

Release date: 2014-12-18

E.58.1. Overview

Major enhancements in PostgreSQL 9.4 include:

• Add jsonb, a more capable and efficient data type for storing JSON data

• Add new SQL command ALTER SYSTEM for changing postgresql.conf configuration file entries

• Reduce lock strength for some ALTER TABLE commands

• Allow materialized views to be refreshed without blocking concurrent reads

• Add support for logical decoding of WAL data, to allow database changes to be streamed out in a customizable format

• Allow background worker processes to be dynamically registered, started and terminated

The above items are explained in more detail in the sections below.

E.58.2. Migration to Version 9.4

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.4 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:
• Tighten checks for multidimensional array input (Bruce Momjian)

Previously, an input array string that started with a single-element sub-array could later contain multi-element sub-arrays, e.g. `'(1), {2,3})'::int[]` would be accepted.

• When converting values of type `date`, `timestamp` or `timestamptz` to JSON, render the values in a format compliant with ISO 8601 (Andrew Dunstan)

Previously such values were rendered according to the current DateStyle setting; but many JSON processors require timestamps to be in ISO 8601 format. If necessary, the previous behavior can be obtained by explicitly casting the datetime value to `text` before passing it to the JSON conversion function.

• The `json #> text[]` path extraction operator now returns its lefthand input, not NULL, if the array is empty (Tom Lane)

This is consistent with the notion that this represents zero applications of the simple field/element extraction operator `->`. Similarly, `json #>> text[]` with an empty array merely coerces its lefthand input to `text`.

• Corner cases in the `JSON` field/element/path extraction operators now return NULL rather than raising an error (Tom Lane)

For example, applying field extraction to a JSON array now yields NULL not an error. This is more consistent (since some comparable cases such as no-such-field already returned NULL), and it makes it safe to create expression indexes that use these operators, since they will now not throw errors for any valid JSON input.

• Cause consecutive whitespace in `to_timestamp()` and `to_date()` format strings to consume a corresponding number of characters in the input string (whitespace or not), then conditionally consume adjacent whitespace, if not in FX mode (Jeevan Chalke)

Previously, consecutive whitespace characters in a non-FX format string behaved like a single whitespace character and consumed all adjacent whitespace in the input string. For example, previous a format string of three spaces would consume only the first space in `' 12'`, but it will now consume all three characters.

• Fix `ts_rank_cd()` to ignore stripped lexemes (Alex Hill)

Previously, stripped lexemes were treated as if they had a default location, producing a rank of dubious usefulness.

• For functions declared to take VARIADIC "any", an actual parameter marked as VARIADIC must be of a determinable array type (Pavel Stehule)

Such parameters can no longer be written as an undecorated string literal or `NULL`; a cast to an appropriate array data type will now be required. Note that this does not affect parameters not marked VARIADIC.

• Ensure that whole-row variables expose the expected column names to functions that pay attention to column names within composite arguments (Tom Lane)

Constructs like `row_to_json(tab.*)` now always emit column names that match the column aliases visible for table `tab` at the point of the call. In previous releases the emitted column names would sometimes be the table's actual column names regardless of any aliases assigned in the query.

• DISCARD now also discards sequence-related state (Fabrízio de Royes Mello, Robert Haas)

• Rename `EXPLAIN ANALYZE`'s "total runtime" output to "execution time" (Tom Lane)

Now that planning time is also reported, the previous name was confusing.
• **SHOW TIME ZONE** now outputs simple numeric UTC offsets in POSIX timezone format (Tom Lane)

  Previously, such timezone settings were displayed as `interval` values. The new output is properly interpreted by `SET TIME ZONE` when passed as a simple string, whereas the old output required special treatment to be re-parsed correctly.

• Foreign data wrappers that support updating foreign tables must consider the possible presence of `AFTER ROW` triggers (Noah Misch)

  When an `AFTER ROW` trigger is present, all columns of the table must be returned by updating actions, since the trigger might inspect any or all of them. Previously, foreign tables never had triggers, so the FDW might optimize away fetching columns not mentioned in the `RETURNING` clause (if any).

• Prevent `CHECK` constraints from referencing system columns, except `tableoid` (Amit Kapila)

  Previously such check constraints were allowed, but they would often cause errors during restores.

• Use the last specified recovery target parameter if multiple target parameters are specified (Heikki Linnakangas)

  Previously, there was an undocumented precedence order among the `recovery_target_` parameters.

• On Windows, automatically preserve quotes in command strings supplied by the user (Heikki Linnakangas)

  User commands that did their own quote preservation might need adjustment. This is likely to be an issue for commands used in `archive_command`, `restore_command`, and `COPY TO/FROM PROGRAM`.

• Remove catalog column `pg_class.reltoastidxid` (Michael Paquier)

• Remove catalog column `pg_rewrite.ev_attr` (Kevin Grittner)

  Per-column rules have not been supported since PostgreSQL 7.3.

• Remove native support for Kerberos authentication (`--with-krb5`, etc) (Magnus Hagander)

  The supported way to use Kerberos authentication is with GSSAPI. The native code has been deprecated since PostgreSQL 8.3.

• In PL/Python, handle domains over arrays like the underlying array type (Rodolfo Campero)

  Previously such values were treated as strings.

• Make `libpq`'s `PQconnectdbParams()` and `PQpingParams()` functions process zero-length strings as defaults (Adrian Vondendriesch)

  Previously, these functions treated zero-length string values as selecting the default in only some cases.

• Change empty arrays returned by the `intarray` module to be zero-dimensional arrays (Bruce Momjian)

  Previously, empty arrays were returned as zero-length one-dimensional arrays, whose text representation looked the same as zero-dimensional arrays (`[]`), but they acted differently in array operations. `intarray`'s behavior in this area now matches the built-in array operators.

• `pg_upgrade` now uses `--username` to specify the user name (Bruce Momjian)
Previously this option was spelled `-u` or `--user`, but that was inconsistent with other tools.

## E.58.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.4 and the previous major release.

### E.58.3.1. Server

- Allow background worker processes to be dynamically registered, started and terminated (Robert Haas)

  The new `worker_spi` module shows an example of use of this feature.

- Allow dynamic allocation of shared memory segments (Robert Haas, Amit Kapila)

  This feature is illustrated in the `test_shmmq` module.

- During crash recovery or immediate shutdown, send uncatchable termination signals (SIGKILL) to child processes that do not shut down promptly (MauMau, Álvaro Herrera)

  This reduces the likelihood of leaving orphaned child processes behind after postmaster shutdown, as well as ensuring that crash recovery can proceed if some child processes have become “stuck”.

- Improve randomness of the database system identifier (Tom Lane)

- Make VACUUM properly report dead but not-yet-removable rows to the statistics collector (Hari Babu)

  Previously these were reported as live rows.

### E.58.3.1.1. Indexes

- Reduce GIN index size (Alexander Korotkov, Heikki Linnakangas)

  Indexes upgraded via `pg_upgrade` will work fine but will still be in the old, larger GIN format. Use `REINDEX` to recreate old GIN indexes in the new format.

- Improve speed of multi-key GIN lookups (Alexander Korotkov, Heikki Linnakangas)

- Add GiST index support for `inet` and `cidr` data types (Emre Hasegeli)

  Such indexes improve subnet and supernet lookups and ordering comparisons.

- Fix rare race condition in B-tree page deletion (Heikki Linnakangas)

- Make the handling of interrupted B-tree page splits more robust (Heikki Linnakangas)

### E.58.3.1.2. General Performance

- Allow multiple backends to insert into WAL buffers concurrently (Heikki Linnakangas)

  This improves parallel write performance.

- Conditionally write only the modified portion of updated rows to WAL (Amit Kapila)

- Improve performance of aggregate functions used as window functions (David Rowley, Florian Pflug, Tom Lane)

- Improve speed of aggregates that use `numeric` state values (Hadi Moshayedi)
• Attempt to freeze tuples when tables are rewritten with CLUSTER or VACUUM FULL (Robert Haas, Andres Freund)

This can avoid the need to freeze the tuples in the future.

• Improve speed of COPY with default nextval() columns (Simon Riggs)

• Improve speed of accessing many different sequences in the same session (David Rowley)

• Raise hard limit on the number of tuples held in memory during sorting and B-tree index builds (Noah Misch)

• Reduce memory allocated by PL/pgSQL DO blocks (Tom Lane)

• Make the planner more aggressive about extracting restriction clauses from mixed AND/OR clauses (Tom Lane)

• Disallow pushing volatile WHERE clauses down into DISTINCT subqueries (Tom Lane)

Pushing down a WHERE clause can produce a more efficient plan overall, but at the cost of evaluating the clause more often than is implied by the text of the query; so don't do it if the clause contains any volatile functions.

• Auto-resize the catalog caches (Heikki Linnakangas)

This reduces memory consumption for sessions accessing only a few tables, and improves performance for sessions accessing many tables.

E.58.3.1.3. Monitoring

• Add pg_stat_archiver system view to report WAL archiver activity (Gabriele Bartolini)

• Add n_mod_since_analyze columns to pg_stat_all_tables and related system views (Mark Kirkwood)

These columns expose the system's estimate of the number of changed tuples since the table's last ANALYZE. This estimate drives decisions about when to auto-analyze.

• Add backend_xid and backend_xmin columns to the system view pg_stat_activity, and a backend_xmin column to pg_stat_replication (Christian Kruse)

E.58.3.1.4. SSL

• Add support for SSL ECDH key exchange (Marko Kreen)

This allows use of Elliptic Curve keys for server authentication. Such keys are faster and have better security than RSA keys. The new configuration parameter ssl_ecdh_curve controls which curve is used for ECDH.

• Improve the default ssl_ciphers setting (Marko Kreen)

• By default, the server not the client now controls the preference order of SSL ciphers (Marko Kreen)

Previously, the order specified by ssl_ciphers was usually ignored in favor of client-side defaults, which are not configurable in most PostgreSQL clients. If desired, the old behavior can be restored via the new configuration parameter ssl_prefer_server_ciphers.

• Make log_connections show SSL encryption information (Andreas Kunert)

• Improve SSL renegotiation handling (Álvaro Herrera)
E.58.3.1.5. Server Settings

- Add new SQL command ALTER SYSTEM for changing `postgresql.conf` configuration file entries (Amit Kapila)
  
  Previously such settings could only be changed by manually editing `postgresql.conf`.

- Add `autovacuum_work_mem` configuration parameter to control the amount of memory used by autovacuum workers (Peter Geoghegan)

- Add `huge_pages` parameter to allow using huge memory pages on Linux (Christian Kruse, Richard Poole, Abhijit Menon-Sen)
  
  This can improve performance on large-memory systems.

- Add `max_worker_processes` parameter to limit the number of background workers (Robert Haas)
  
  This is helpful in configuring a standby server to have the required number of worker processes (the same as the primary).

- Add `superuser-only session_preload_libraries` parameter to load libraries at session start (Peter Eisentraut)
  
  In contrast to `local_preload_libraries`, this parameter can load any shared library, not just those in the `$libdir/plugins` directory.

- Add `wal_log_hints` parameter to enable WAL logging of hint-bit changes (Sawada Masahiko)
  
  Hint bit changes are not normally logged, except when checksums are enabled. This is useful for external tools like `pg_rewind`.

- Increase the default settings of `work_mem` and `maintenance_work_mem` by four times (Bruce Momjian)
  
  The new defaults are 4MB and 64MB respectively.

- Increase the default setting of `effective_cache_size` to 4GB (Bruce Momjian, Tom Lane)

- Allow `printf`-style space padding to be specified in `log_line_prefix` (David Rowley)

- Allow terabyte units (TB) to be used when specifying configuration variable values (Simon Riggs)

- Show PIDs of lock holders and waiters and improve information about relations in `log_lock_waits` log messages (Christian Kruse)

- Reduce server logging level when loading shared libraries (Peter Geoghegan)
  
  The previous level was `LOG`, which was too verbose for libraries loaded per-session.

- On Windows, make `SQL_ASCII`-encoded databases and server processes (e.g., `postmaster`) emit messages in the character encoding of the server's Windows user locale (Alexander Law, Noah Misch)
  
  Previously these messages were output in the Windows ANSI code page.

E.58.3.2. Replication and Recovery

- Add replication slots to coordinate activity on streaming standbys with the node they are streaming from (Andres Freund, Robert Haas)
  
  Replication slots allow preservation of resources like WAL files on the primary until they are no longer needed by standby servers.
• Add recovery parameter recovery_min_apply_delay to delay replication (Robert Haas, Fabrício de Royes Mello, Simon Riggs)

Delaying replay on standby servers can be useful for recovering from user errors.

• Add recovery target option immediate to stop WAL recovery as soon as a consistent state is reached (MauMau, Heikki Linnakangas)

• Improve recovery target processing (Heikki Linnakangas)

The timestamp reported by pg_last_xact_replay_timestamp() now reflects already-committed records, not transactions about to be committed. Recovering to a restore point now replays the restore point, rather than stopping just before the restore point.

• pg_switch_xlog() now clears any unused trailing space in the old WAL file (Heikki Linnakangas)

This improves the compression ratio for WAL files.

• Report failure return codes from external recovery commands (Peter Eisentraut)

• Reduce spinlock contention during WAL replay (Heikki Linnakangas)

• Write WAL records of running transactions more frequently (Andres Freund)

This allows standby servers to start faster and clean up resources more aggressively.

E.58.3.2.1. Logical Decoding

Logical decoding allows database changes to be streamed in a configurable format. The data is read from the WAL and transformed into the desired target format. To implement this feature, the following changes were made:

• Add support for logical decoding of WAL data, to allow database changes to be streamed out in a customizable format (Andres Freund)

• Add new wal_level setting logical to enable logical change-set encoding in WAL (Andres Freund)

• Add table-level parameter REPLICA IDENTITY to control logical replication (Andres Freund)

• Add relation option user_catalog_table to identify user-created tables involved in logical change-set encoding (Andres Freund)

• Add pg_recvlogical application to receive logical-decoding data (Andres Freund)

• Add test_decoding module to illustrate logical decoding at the SQL level (Andres Freund)

E.58.3.3. Queries

• Add WITH ORDINALITY syntax to number the rows returned from a set-returning function in the FROM clause (Andrew Gierth, David Fetter)

  This is particularly useful for functions like unnest().

• Add ROWS FROM() syntax to allow horizontal concatenation of set-returning functions in the FROM clause (Andrew Gierth)

• Allow SELECT to have an empty target list (Tom Lane)

  This was added so that views that select from a table with zero columns can be dumped and restored correctly.
• Ensure that SELECT ... FOR UPDATE NOWAIT does not wait in corner cases involving already-concurrently-updated tuples (Craig Ringer and Thomas Munro)

**E.58.3.4. Utility Commands**

• Add DISCARD SEQUENCES command to discard cached sequence-related state (Fabrizio de Royes Mello, Robert Haas)

  DISCARD ALL will now also discard such information.

• Add FORCE NULL option to COPY FROM, which causes quoted strings matching the specified null string to be converted to NULLs in CSV mode (Ian Barwick, Michael Paquier)

  Without this option, only unquoted matching strings will be imported as null values.

• Issue warnings for commands used outside of transaction blocks when they can have no effect (Bruce Momjian)

  New warnings are issued for SET LOCAL, SET CONSTRAINTS, SET TRANSACTION and ABORT when used outside a transaction block.

**E.58.3.4.1. EXPLAIN**

• Make EXPLAIN ANALYZE show planning time (Andreas Karlsson)

• Make EXPLAIN show the grouping columns in Agg and Group nodes (Tom Lane)

• Make EXPLAIN ANALYZE show exact and lossy block counts in bitmap heap scans (Etsuro Fujita)

**E.58.3.4.2. Views**

• Allow a materialized view to be refreshed without blocking other sessions from reading the view meanwhile (Kevin Grittner)

  This is done with REFRESH MATERIALIZED VIEW CONCURRENTLY.

• Allow views to be automatically updated even if they contain some non-updatable columns (Dean Rasheed)

  Previously the presence of non-updatable output columns such as expressions, literals, and function calls prevented automatic updates. Now INSERTs, UPDATEs and DELETEs are supported, provided that they do not attempt to assign new values to any of the non-updatable columns.

• Allow control over whether INSERTs and UPDATEs can add rows to an auto-updatable view that would not appear in the view (Dean Rasheed)

  This is controlled with the new CREATE VIEW clause WITH CHECK OPTION.

• Allow security barrier views to be automatically updatable (Dean Rasheed)

**E.58.3.5. Object Manipulation**

• Support triggers on foreign tables (Ronan Dunklau)

• Allow moving groups of objects from one tablespace to another using the ALL IN TABLESPACE ... SET TABLESPACE form of ALTER TABLE, ALTER INDEX, or ALTER MATERIALIZED VIEW (Stephen Frost)

• Allow changing foreign key constraint deferrability via ALTER TABLE ... ALTER CONSTRAINT (Simon Riggs)
• Reduce lock strength for some ALTER TABLE commands (Simon Riggs, Noah Misch, Robert Haas)

  Specifically, VALIDATE CONSTRAINT, CLUSTER ON, SET WITHOUT CLUSTER, ALTER COLUMN SET STATISTICS, ALTER COLUMN SET (attribute_option), ALTER COLUMN RESET (attribute_option) no longer require ACCESS EXCLUSIVE locks.

• Allow tablespace options to be set in CREATE TABLESPACE (Vik Fearing)

  Formerly these options could only be set via ALTER TABLESPACE.

• Allow CREATE AGGREGATE to define the estimated size of the aggregate's transition state data (Hadi Moshayedi)

  Proper use of this feature allows the planner to better estimate how much memory will be used by aggregates.

• Fix DROP IF EXISTS to avoid errors for non-existent objects in more cases (Pavel Stehule, Dean Rasheed)

• Improve how system relations are identified (Andres Freund, Robert Haas)

  Previously, relations once moved into the pg_catalog schema could no longer be modified or dropped.

E.58.3.6. Data Types

• Fully implement the line data type (Peter Eisentraut)

  The line segment data type (lseg) has always been fully supported. The previous line data type (which was enabled only via a compile-time option) is not binary or dump-compatible with the new implementation.

• Add pg_lsn data type to represent a WAL log sequence number (LSN) (Robert Haas, Michael Paquier)

• Allow single-point polygons to be converted to circles (Bruce Momjian)

• Support time zone abbreviations that change UTC offset from time to time (Tom Lane)

  Previously, PostgreSQL assumed that the UTC offset associated with a time zone abbreviation (such as EST) never changes in the usage of any particular locale. However this assumption fails in the real world, so introduce the ability for a zone abbreviation to represent a UTC offset that sometimes changes. Update the zone abbreviation definition files to make use of this feature in timezone locales that have changed the UTC offset of their abbreviations since 1970 (according to the IANA timezone database). In such timezones, PostgreSQL will now associate the correct UTC offset with the abbreviation depending on the given date.

• Allow 5+ digit years for non-ISO timestamp and date strings, where appropriate (Bruce Momjian)

• Add checks for overflow/underflow of interval values (Bruce Momjian)

E.58.3.6.1. JSON

• Add jsonb, a more capable and efficient data type for storing JSON data (Oleg Bartunov, Teodor Sigaev, Alexander Korotkov, Peter Geoghegan, Andrew Dunstan)

  This new type allows faster access to values within a JSON document, and faster and more useful indexing of JSON columns. Scalar values in jsonb documents are stored as appropriate scalar
SQL types, and the JSON document structure is pre-parsed rather than being stored as text as in the original json data type.

- Add new JSON functions to allow for the construction of arbitrarily complex JSON trees (Andrew Dunstan, Laurence Rowe)

  New functions include `json_array_elements_text()`, `json_build_array()`, `json_object()`, `json_object_agg()`, `json_to_record()`, and `json_to_recordset()`.

- Add `json_typeof()` to return the data type of a json value (Andrew Tipton)

### E.58.3.7. Functions

- Add `pg_sleep_for(interval)` and `pg_sleep_until(timestamp)` to specify delays more flexibly (Vik Fearing, Julien Rouhaud)

  The existing `pg_sleep()` function only supports delays specified in seconds.

- Add `cardinality()` function for arrays (Marko Tiikkaja)

  This returns the total number of elements in the array, or zero for an array with no elements.

- Add SQL functions to allow large object reads/writes at arbitrary offsets (Pavel Stehule)

- Allow `unnest()` to take multiple arguments, which are individually unnested then horizontally concatenated (Andrew Gierth)

- Add functions to construct `times`, `dates`, `timestamps`, `timetz`, and `intervals` from individual values, rather than strings (Pavel Stehule)

  These functions' names are prefixed with `make_`, e.g. `make_date()`.

- Make `to_char()`’s `TZ` format specifier return a useful value for simple numeric time zone offsets (Tom Lane)

  Previously, `to_char(CURRENT_TIMESTAMP, 'TZ')` returned an empty string if the time zone was set to a constant like `-4`.

- Add timezone offset format specifier `OF` to `to_char()` (Bruce Momjian)

- Improve the random seed used for `random()` (Honza Horak)

- Tighten validity checking for Unicode code points in `chr(int)` (Tom Lane)

  This function now only accepts values that are valid UTF8 characters according to RFC 3629.

### E.58.3.7.1. System Information Functions

- Add functions for looking up objects in `pg_class`, `pg_proc`, `pg_type`, and `pg_operator` that do not generate errors for non-existent objects (Yugo Nagata, Nozomi Anzai, Robert Haas)

  For example, `to_regclass()` does a lookup in `pg_class` similarly to the `regclass` input function, but it returns NULL for a non-existent object instead of failing.

- Add function `pg_filenode_relation()` to allow for more efficient lookup of relation names from filenodes (Andres Freund)

- Add `parameter_default` column to `information_schema.parameters` view (Peter Eisentraut)

- Make `information_schema.schemata` show all accessible schemas (Peter Eisentraut)
Previously it only showed schemas owned by the current user.

**E.58.3.7.2. Aggregates**

- Add control over which rows are passed into aggregate functions via the `FILTER` clause (David Fetter)
- Support ordered-set (`WITHIN GROUP`) aggregates (Atri Sharma, Andrew Gierth, Tom Lane)
- Add standard ordered-set aggregates `percentile_cont()`, `percentile_disc()`, `mode()`, `rank()`, `dense_rank()`, `percent_rank()`, and `cume_dist()` (Atri Sharma, Andrew Gierth)
- Support `VARIADIC` aggregate functions (Tom Lane)
- Allow polymorphic aggregates to have non-polymorphic state data types (Tom Lane)

This allows proper declaration in SQL of aggregates like the built-in aggregate `array_agg()`.

**E.58.3.8. Server-Side Languages**

- Add event trigger support to PL/Perl and PL/Tcl (Dimitri Fontaine)
- Convert numeric values to decimal in PL/Python (Szymon Guz, Ronan Dunklau)

Previously such values were converted to Python `float` values, risking loss of precision.

**E.58.3.8.1. PL/pgSQL Server-Side Language**

- Add ability to retrieve the current PL/pgSQL call stack using `GET DIAGNOSTICS` (Pavel Stehule, Stephen Frost)
- Add option `print_strict_params` to display the parameters passed to a query that violated a `STRICT` constraint (Marko Tiikkaja)
- Add variables `plpgsql.extra_warnings` and `plpgsql.extra_errors` to enable additional PL/pgSQL warnings and errors (Marko Tiikkaja, Petr Jelinek)

Currently only warnings/errors about shadowed variables are available.

**E.58.3.9. libpq**

- Make `libpq`'s `PQconndefaults()` function ignore invalid service files (Steve Singer, Bruce Momjian)

Previously it returned `NULL` if an incorrect service file was encountered.
- Accept TLS protocol versions beyond `TLSv1` in `libpq` (Marko Kreen)

**E.58.3.10. Client Applications**

- Add `createuser` option `-g` to specify role membership (Christopher Browne)
- Add `vaccumdb` option `--analyze-in-stages` to analyze in stages of increasing granularity (Peter Eisentraut)

This allows minimal statistics to be created quickly.
- Make `pg_resettxlog` with option `-n` output current and potentially changed values (Rajeev Rastogi)
• Make initdb throw error for incorrect locale settings, rather than silently falling back to a default choice (Tom Lane)

• Make pg_ctl return exit code 4 for an inaccessible data directory (Amit Kapila, Bruce Momjian)
   This behavior more closely matches the Linux Standard Base (LSB) Core Specification.

• On Windows, ensure that a non-absolute -D path specification is interpreted relative to pg_ctl's current directory (Kumar Rajeev Rastogi)
   Previously it would be interpreted relative to whichever directory the underlying Windows service was started in.

• Allow sizeof() in ECPG C array definitions (Michael Meskes)

• Make ECPG properly handle nesting of C-style comments in both C and SQL text (Michael Meskes)

E.58.3.10.1. psql

• Suppress "No rows" output in psql expanded mode when the footer is disabled (Bruce Momjian)

• Allow Control-C to abort psql when it's hung at connection startup (Peter Eisentraut)

E.58.3.10.1.1. Backslash Commands

• Make psql's \db+ show tablespace options (Magnus Hagander)

• Make \do+ display the functions that implement the operators (Marko Tiikkaja)

• Make \d+ output an OID line only if an oid column exists in the table (Bruce Momjian)
   Previously, the presence or absence of an oid column was always reported.

• Make \d show disabled system triggers (Bruce Momjian)
   Previously, if you disabled all triggers, only user triggers would show as disabled.

• Fix \copy to no longer require a space between stdin and a semicolon (Etsuro Fujita)

• Output the row count at the end of \copy, just like COPY already did (Kumar Rajeev Rastogi)

• Fix \conninfo to display the server's IP address for connections using hostaddr (Fujii Masao)
   Previously \conninfo could not display the server's IP address in such cases.

• Show the SSL protocol version in \conninfo (Marko Kreen)

• Add tab completion for \pset (Pavel Stehule)

• Allow \pset with no arguments to show all settings (Gilles Darold)

• Make \s display the name of the history file it wrote without converting it to an absolute path (Tom Lane)
   The code previously attempted to convert a relative file name to an absolute path for display, but frequently got it wrong.

E.58.3.10.2. pg_dump

• Allow pg_restore options -I, -P, -T and -n to be specified multiple times (Heikki Linnakangas)
   This allows multiple objects to be restored in one operation.
• Optionally add `IF EXISTS` clauses to the `DROP` commands emitted when removing old objects during a restore (Pavel Stehule)

This change prevents unnecessary errors when removing old objects. The new `--if-exists` option for `pg_dump`, `pg_dumpall`, and `pg_restore` is only available when `--clean` is also specified.

**E.58.3.10.3. pg_basebackup**

• Add `pg_basebackup` option `--xlogdir` to specify the `pg_xlog` directory location (Haribabu Kommi)

• Allow `pg_basebackup` to relocate tablespaces in the backup copy (Steeve Lennmark)

This is particularly useful for using `pg_basebackup` on the same machine as the primary.

• Allow network-stream base backups to be throttled (Antonin Houska)

This can be controlled with the `pg_basebackup` `--max-rate` parameter.

**E.58.3.11. Source Code**

• Improve the way tuples are frozen to preserve forensic information (Robert Haas, Andres Freund)

This change removes the main objection to freezing tuples as soon as possible. Code that inspects tuple flag bits will need to be modified.

• No longer require function prototypes for functions marked with the `PG_FUNCTION_INFO_V1` macro (Peter Eisentraut)

This change eliminates the need to write boilerplate prototypes. Note that the `PG_FUNCTION_INFO_V1` macro must appear before the corresponding function definition to avoid compiler warnings.

• Remove `SnapshotNow` and `HeapTupleSatisfiesNow()` (Robert Haas)

All existing uses have been switched to more appropriate snapshot types. Catalog scans now use MVCC snapshots.

• Add an API to allow memory allocations over one gigabyte (Noah Misch)

• Add `psprintf()` to simplify memory allocation during string composition (Peter Eisentraut, Tom Lane)

• Support `printf()` size modifier `z` to print `size_t` values (Andres Freund)

• Change API of `appendStringInfoVA()` to better use `vsnprintf()` (David Rowley, Tom Lane)

• Allow new types of external toast datums to be created (Andres Freund)

• Add single-reader, single-writer, lightweight shared message queue (Robert Haas)

• Improve spinlock speed on x86_64 CPUs (Heikki Linnakangas)

• Remove spinlock support for unsupported platforms SINIX, Sun3, and NS32K (Robert Haas)

• Remove IRIX port (Robert Haas)

• Reduce the number of semaphores required by `--disable-spinlocks` builds (Robert Haas)

• Rewrite `duplicate_oids` Unix shell script in Perl (Andrew Dunstan)

• Add Test Anything Protocol (TAP) tests for client programs (Peter Eisentraut)
Currently, these tests are run by `make check-world` only if the `--enable-tap-tests` option was given to configure. This might become the default behavior in some future release.

- Add make targets `check-tests` and `installcheck-tests`, which allow selection of individual tests to be run (Andrew Dunstan)

- Remove `maintainer-check` makefile rule (Peter Eisentraut)

  The default build rules now include all the formerly-optional tests.

- Improve support for `VPATH` builds of PGXS modules (Cédric Villemain, Andrew Dunstan, Peter Eisentraut)

- Upgrade to Autoconf 2.69 (Peter Eisentraut)

- Add a configure flag that appends custom text to the `PG_VERSION` string (Oskari Saarenmaa)

  This is useful for packagers building custom binaries.

- Improve DocBook XML validity (Peter Eisentraut)

- Fix various minor security and sanity issues reported by the Coverity scanner (Stephen Frost)

- Improve detection of invalid memory usage when testing PostgreSQL with Valgrind (Noah Misch)

- Improve sample Emacs configuration file `emacs.samples` (Peter Eisentraut)

  Also add `.dir-locals.el` to the top of the source tree.

- Allow `pgindent` to accept a command-line list of typedefs (Bruce Momjian)

- Make `pgindent` smarter about blank lines around preprocessor conditionals (Bruce Momjian)

- Avoid most uses of `dlltool` in Cygwin and Mingw builds (Marco Atzeri, Hiroshi Inoue)

- Support client-only installs in MSVC (Windows) builds (MauMau)

### E.58.3.12. Additional Modules

- Add `pg_prewarm` extension to preload relation data into the shared buffer cache at server start (Robert Haas)

  This allows reaching full operating performance more quickly.

- Add UUID random number generator `gen_random_uuid()` to `pgcrypto` (Oskari Saarenmaa)

  This allows creation of version 4 UUIDs without requiring installation of uuid-ossp.

- Allow `uuid-ossp` to work with the BSD or e2fsprogs UUID libraries, not only the OSSP UUID library (Matteo Beccati)

  This improves the `uuid-ossp` module's portability since it no longer has to have the increasingly-obsolete OSSP library. The module's name is now rather a misnomer, but we won't change it.

- Add option to `auto_explain` to include trigger execution time (Horiguchi Kyotaro)

- Fix `pgstattuple` to not report rows from uncommitted transactions as dead (Robert Haas)

- Make `pgstattuple` functions use `regclass-type` arguments (Satoshi Nagayasu)

  While `text-type` arguments are still supported, they may be removed in a future major release.
• Improve consistency of pgrowlocks output to honor snapshot rules more consistently (Robert Haas)
• Improve pg_trgm's choice of trigrams for indexed regular expression searches (Alexander Korotkov)

  This change discourages use of trigrams containing whitespace, which are usually less selective.
• Allow pg_xlogdump to report a live log stream with --follow (Heikki Linnakangas)
• Store cube data more compactly (Stas Kelvich)

  Existing data must be dumped/restored to use the new format. The old format can still be read.
• Reduce vacuumlo client-side memory usage by using a cursor (Andrew Dunstan)
• Dramatically reduce memory consumption in pg_upgrade (Bruce Momjian)
• Pass pg_upgrade's user name (--U) option to generated analyze scripts (Bruce Momjian)

E.58.12.1. pgbench

• Remove line length limit for pgbench scripts (Sawada Masahiko)

  The previous line limit was BUFSIZ.
• Add long option names to pgbench (Fabien Coelho)
• Add pgbench option --rate to control the transaction rate (Fabien Coelho)
• Add pgbench option --progress to print periodic progress reports (Fabien Coelho)

E.58.12.2. pg_stat_statements

• Make pg_stat_statements use a file, rather than shared memory, for query text storage (Peter Geoghegan)

  This removes the previous limitation on query text length, and allows a higher number of unique statements to be tracked by default.
• Allow reporting of pg_stat_statements's internal query hash identifier (Daniel Farina, Sameer Thakur, Peter Geoghegan)
• Add the ability to retrieve all pg_stat_statements information except the query text (Peter Geoghegan)

  This allows monitoring tools to fetch query text only for just-created entries, improving performance during repeated querying of the statistics.
• Make pg_stat_statements ignore DEALLOCATE commands (Fabien Coelho)

  It already ignored PREPARE, as well as planning time in general, so this seems more consistent.
• Save the statistics file into $PGDATA/pg_stat at server shutdown, rather than $PGDATA/TABLE/global (Fujii Masao)

E.59. Release 9.3.25

Release date: 2018-11-08

This release contains a variety of fixes from 9.3.24. For information about new features in the 9.3 major release, see Section E.84.
This is expected to be the last PostgreSQL release in the 9.3.X series. Users are encouraged to update to a newer release branch soon.

**E.59.1. Migration to Version 9.3.25**

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.23, see Section E.61.

**E.59.2. Changes**

- **Fix corner-case failures in has_foo_privilege() family of functions (Tom Lane)**
  
  Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. has_column_privilege() was additionally capable of crashing on some platforms.

- **Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)**

- **Avoid O(N^3) slowdown in lexer for long strings of + or - characters (Andrew Gierth)**

- **Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)**

- **Fix failure of UPDATE/DELETE ... WHERE CURRENT OF ... after rewinding the referenced cursor (Tom Lane)**
  
  A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

- **Fix EvalPlanQual to handle conditionally-executed InitPlans properly (Andrew Gierth, Tom Lane)**
  
  This resulted in hard-to-reproduce crashes or wrong answers in concurrent updates, if they contained code such as an uncorrelated sub-SELECT inside a CASE construct.

- **Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF (Tom Lane, Kenji Uno)**
  
  This bug affected full-text-search operations, as well as contrib/ltree and contrib/pg_trgm.

- **Ensure that sequences owned by a foreign table are processed by ALTER OWNER on the table (Peter Eisentraut)**
  
  The ownership change should propagate to such sequences as well, but this was missed for foreign tables.

- **Fix over-allocation of space for array_out()’s result string (Keiichi Hirobe)**

- **Fix memory leak in repeated SP-GiST index scans (Tom Lane)**
  
  This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- **Avoid crash if a utility command causes infinite recursion (Tom Lane)**

- **When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master (Michael Paquier, Konstantin Knizhnik)**

- **Randomize the random() seed in bootstrap and standalone backends, and in initdb (Noah Misch)**
The main practical effect of this change is that it avoids a scenario where initdb might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

- Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)
  This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.
- Don't run atexit callbacks when servicing SIGQUIT (Heikki Linnakangas)
- Don't record foreign-server user mappings as members of extensions (Tom Lane)
  If CREATE USER MAPPING is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can't be extension members, so user mappings shouldn't be either.
- Make syslogger more robust against failures in opening CSV log files (Tom Lane)
- Fix possible inconsistency in pg_dump's sorting of dissimilar object names (Jacob Champion)
- Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)
  This avoids failures due to the new policy of running restores with restrictive search path.
- Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)
  pg_upgrade failed to preserve an event trigger's extension-membership status.
- Fix pg_upgrade's cluster state check to work correctly on a standby server (Bruce Momjian)
- Enforce type cube's dimension limit in all contrib/cube functions (Andrey Borodin)
  Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.
- Fix contrib/unaccent's unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)
  Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.
- Fix build problems on macOS 10.14 (Mojave) (Tom Lane)
  Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSRoot variable in the arguments of configure or make.
  It is now recommended that Perl-related extensions write $(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.
  Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.
- Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)
  Perl no longer includes the current directory in its search path by default; work around that.
• Support building on Windows with Visual Studio 2015 or Visual Studio 2017 (Michael Paquier, Haribabu Kommi)

• Allow btree comparison functions to return INT_MIN (Tom Lane)

Up to now, we’ve forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of memcmp(), strcmp(), etc, as POSIX doesn’t place any such restriction on those functions. At least some recent versions of memcmp() can return INT_MIN, causing incorrect sort ordering. Hence, we’ve removed this restriction. Callers must now use the INVERT_COMPARE_RESULT() macro if they wish to invert the sort order.

• Fix recursion hazard in shared-invalidation message processing (Tom Lane)

This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.

This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.

• Save and restore SPI’s global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)

This prevents possible interference when one SPI-using function calls another.

• Provide ALLOCSET_DEFAULT_SIZES and sibling macros in back branches (Tom Lane)

These macros have existed since 9.6, but there were requests to add them to older branches to allow extensions to rely on them without branch-specific coding.

• Avoid using potentially-under-aligned page buffers (Tom Lane)

Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can’t place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.

• Make src/port/snprintf.c follow the C99 standard’s definition of snprintf()’s result value (Tom Lane)

On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.

• When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)

This avoids problems with missed floating point overflow checks.

• Fix configure’s detection of the result type of strerror_r() (Tom Lane)

The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.

• Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.60. Release 9.3.24

Release date: 2018-08-09
This release contains a variety of fixes from 9.3.23. For information about new features in the 9.3 major release, see Section E.84.

The PostgreSQL community will stop releasing updates for the 9.3.X release series shortly after September 2018. Users are encouraged to update to a newer release branch soon.

E.60.1. Migration to Version 9.3.24

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.23, see Section E.61.

E.60.2. Changes

• Fix failure to reset libpq's state fully between connection attempts (Tom Lane)

An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general, an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasichkov for reporting this issue. (CVE-2018-10915)

• Ensure that updates to the relfrozenxid and relminmxid values for “nailed” system catalogs are processed in a timely fashion (Andres Freund)

Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as pg_authid, because the stale cache data could persist into new sessions as well as existing ones.

• Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

• Avoid emitting a bogus WAL record when recycling an all-zero btree page (Amit Kapila)

This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

• Improve performance of WAL replay for transactions that drop many relations (Fujii Masao)

This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

• Improve performance of lock releasing in standby server WAL replay (Thomas Munro)

• Ensure a table's cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)

Previously, the failed index's OID could remain in the list, causing problems later in the same session.

• Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)
This resulted in failure to recognize that an index on a composite column could provide the sort order needed for a mergejoin on that column.

- Fix SQL-standard `FETCH FIRST` syntax to allow parameters ($n), as the standard expects (Andrew Gierth)
- Fix failure to schema-qualify some object names in `getObjectDescription` output (Kyotaro Horiguchi, Tom Lane)

Names of collations, conversions, and text search objects were not schema-qualified when they should be.

- Widen `COPY FROM`'s current-line-number counter from 32 to 64 bits (David Rowley)
  This avoids two problems with input exceeding 4G lines: `COPY FROM WITH HEADER` would drop a line every 4G lines, not only the first line, and error reports could show a wrong line number.

- Add a string freeing function to ecpg's `pgtypes` library, so that cross-module memory management problems can be avoided on Windows (Takayuki Tsunakawa)
  On Windows, crashes can ensue if the `free` call for a given chunk of memory is not made from the same DLL that `malloc`ed the memory. The `pgtypes` library sometimes returns strings that it expects the caller to free, making it impossible to follow this rule. Add a `PGTYPESchar_free()` function that just wraps `free`, allowing applications to follow this rule.

- Fix ecpg's support for `long long` variables on Windows, as well as other platforms that declare `strtoll/strtoull` nonstandardly or not at all (Dang Minh Huong, Tom Lane)

- Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in the semantics of a statement intra-session (Tom Lane)
  This error led to assertion failures, or in rare cases, failure to enforce the `INTO STRICT` option as expected.

- Fix password prompting in client programs so that echo is properly disabled on Windows when `stdin` is not the terminal (Matthew Stickney)

- Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)
  The previous fix for quoting of `search_path` and other list-valued variables in `pg_dump` output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

- Make `pg_upgrade` check that the old server was shut down cleanly (Bruce Momjian)
  The previous check could be fooled by an immediate-mode shutdown.

- Fix crash in `contrib/ltrie/s.lca()` function when the input array is empty (Pierre Ducroquet)

- Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

- Rearrange makefiles to ensure that programs link to freshly-built libraries (such as `libpq.so`) rather than ones that might exist in the system library directories (Tom Lane)
  This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

- Update time zone data files to `tzdata` release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.
  This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard
time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbreviations in use; the only known effect is that the is_dst column in the pg_timezone_names view will now be true in winter and false in summer in these cases.

E.61. Release 9.3.23

Release date: 2018-05-10

This release contains a variety of fixes from 9.3.22. For information about new features in the 9.3 major release, see Section E.84.

E.61.1. Migration to Version 9.3.23

A dump/restore is not required for those running 9.3.X.

However, if the function marking mistakes mentioned in the first changelog entry below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 9.3.22, see Section E.62.

E.61.2. Changes

• Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

  The functions query_to_xml, cursor_to_xml, cursor_to_xmlschema, query_to_xmlschema, and query_to_xml_and_xmlschema should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions’ pg_proc entries, for example ALTER FUNCTION pg_catalog.query_to_xml(text, boolean, boolean, text) VOLATILE. (Note that that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

• Avoid re-using TOAST value OIDS that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

  Once the OID counter has wrapped around, it’s possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value nnnnn” errors, which would persist until the dead entry was removed by VACUUM. Fix by not selecting such OIDs when creating a new TOAST entry.

• Change ANALYZE’s algorithm for updating pg_class.reltuples (David Gould)

  Previously, pages not actually scanned by ANALYZE were assumed to retain their old tuple density. In a large table where ANALYZE samples only a small fraction of the pages, this meant that the overall tuple density estimate could not change very much, so that reltuples would change nearly proportionally to changes in the table’s physical size (relpages) regardless of what was actually happening in the table. This has been observed to result in reltuples becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that ANALYZE’s sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

• Fix UPDATE/DELETE ... WHERE CURRENT OF to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)

• Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)
This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.

• Fix misoptimization of CHECK constraints having provably-NULL subclauses of top-level AND/OR conditions (Tom Lane, Dean Rasheed)

This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.

• Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)

• Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)

• Fix overly strict sanity check in heap_prepare_freeze_tuple (Álvaro Herrera)

This could result in incorrect “cannot freeze committed xmax” failures in databases that have been pg_upgrade'd from 9.2 or earlier.

• Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)

• Reduce locking during autovacuum worker scheduling (Jeff Janes)

The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.

• Ensure client hostname is copied while copying pg_stat_activity data to local memory (Edmund Horner)

Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the client hostname column to change unexpectedly if any existing session disconnected.

• Fix incorrect processing of multiple compound affixes in ispell dictionaries (Arthur Zakirov)

• Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on text columns (Tom Lane)

Such searches would return the wrong set of rows in most non-C locales.

• Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)

Previously, the tuple count was reported to be the same as that of the underlying table, which is wrong if the index is partial.

• Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)

Previously it reported the estimated number of heap tuples, which might be inaccurate, and is certainly wrong if the index is partial.

• Allow scalarltsel and scalargtsel to be used on non-core datatypes (Tomas Vondra)

• Reduce libpq's memory consumption when a server error is reported after a large amount of query output has been collected (Tom Lane)

Discard the previous output before, not after, processing the error message. On some platforms, notably Linux, this can make a difference in the application's subsequent memory footprint.

• Fix double-free crashes in ecpg (Patrick Krecker, Jeevan Ladhe)

• Fix ecpg to handle long long int variables correctly in MSVC builds (Michael Meskes, Andrew Gierth)
• Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)

The local_preload_libraries, session_preload_libraries, shared_preload_libraries, and temp_tablespaces variables were not correctly quoted in pg_dump output. This would cause problems if settings for these variables appeared in CREATE FUNCTION ... SET or ALTER DATABASE/ROLE ... SET clauses.

• Fix overflow handling in PL/pgSQL integer FOR loops (Tom Lane)

The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

• Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisentraut)

• Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)

• Rename internal b64_encode and b64_decode functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)

• Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)

This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.

• Update time zone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

E.62. Release 9.3.22

Release date: 2018-03-01

This release contains a variety of fixes from 9.3.21. For information about new features in the 9.3 major release, see Section E.84.

E.62.1. Migration to Version 9.3.22

A dump/restore is not required for those running 9.3.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that you read the documentation changes described in the first changelog entry below, and take suitable steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 9.3.18, see Section E.66.

E.62.2. Changes

• Document how to configure installations and applications to guard against search-path-dependent trojan-horse attacks from other users (Noah Misch)

Using a search_path setting that includes any schemas writable by a hostile user enables that user to capture control of queries and then run arbitrary SQL code with the permissions of the attacked user. While it is possible to write queries that are proof against such hijacking, it is notation-
ally tedious, and it's very easy to overlook holes. Therefore, we now recommend configurations in which no untrusted schemas appear in one's search path. Relevant documentation appears in Section 5.8.6 (for database administrators and users), Section 34.1 (for application authors), Section 38.16.1 (for extension authors), and CREATE FUNCTION (for authors of SECURITY DEFINER functions). (CVE-2018-1058)

- Avoid use of insecure search_path settings in pg_dump and other client programs (Noah Misch, Tom Lane)

pg_dump, pg_upgrade, vacuumdb and other PostgreSQL-provided applications were themselves vulnerable to the type of hijacking described in the previous changelog entry; since these applications are commonly run by superusers, they present particularly attractive targets. To make them secure whether or not the installation as a whole has been secured, modify them to include only the pg_catalog schema in their search_path settings. Autovacuum worker processes now do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example, user-provided functions in index expressions — the tighter search_path may result in errors, which will need to be corrected by adjusting those user-provided functions to not assume anything about what search path they are invoked under. That has always been good practice, but now it will be necessary for correct behavior. (CVE-2018-1058)

- Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom Lane)

If a CTE (WITH clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

- Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do not match mergeclauses” planner errors in corner cases.

- Repair pg_upgrade's failure to preserve relfrozenxid for materialized views (Tom Lane, Andres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The problem would be more likely to occur in seldom-refreshed materialized views, or ones that were maintained only with REFRESH MATERIALIZED VIEW CONCURRENTLY.

If such corruption is observed, it can be repaired by refreshing the materialized view (without CONCURRENTLY).

- Fix incorrect reporting of PL/Python function names in error CONTEXT stacks (Tom Lane)

An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function’s name twice, rather than the expected results. Also, an error in a nested PL/Python DO block could result in a null pointer dereference crash on some platforms.

- Allow contrib/auto_explain's log_min_duration setting to range up to INT_MAX, or about 24 days instead of 35 minutes (Tom Lane)

E.63. Release 9.3.21

Release date: 2018-02-08

This release contains a variety of fixes from 9.3.20. For information about new features in the 9.3 major release, see Section E.84.
**E.63.1. Migration to Version 9.3.21**

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.18, see Section E.66.

**E.63.2. Changes**

- Ensure that all temporary files made by pg_upgrade are non-world-readable (Tom Lane, Noah Misch)

  pg_upgrade normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing `pg_dumpall -g` output would be group- or world-readable, or even writable, if the user's `umask` setting allows. In typical usage on multi-user machines, the `umask` and/or the working directory’s permissions would be tight enough to prevent problems; but there may be people using pg_upgrade in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

- Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

  In some cases `VACUUM` would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

- Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)

  These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.

- Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by `ALTER TABLE NO INHERIT` (Tom Lane)

- Repair failure with correlated sub-`SELECT` inside `VALUES` inside a LATERAL subquery (Tom Lane)

- Fix “could not devise a query plan for the given query” planner failure for some cases involving nested `UNION ALL` inside a lateral subquery (Tom Lane)

- Fix `has_sequence_privilege()` to support `WITH GRANT OPTION` tests, as other privilege-testing functions do (Joe Conway)

- In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)

  We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don’t promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only `xpath()` and related functions; other XML code paths already acted this way.

- Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)

  Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.

- Prevent stack-overflow crashes when planning extremely deeply nested set operations (`UNION/INTERSECT/EXCEPT`) (Tom Lane)
• Fix null-pointer crashes for some types of LDAP URLs appearing in pg_hba.conf (Thomas Munro)

• Fix sample INSTR() functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)

These functions are stated to be Oracle® compatible, but they weren't exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.

The sample code has been adjusted to match Oracle's behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

• Fix pg_dump to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with DATABASE so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with LARGE_OBJECT. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.

Note that this change has user-visible results in the output of pg_restore --list.

• In ecpg, detect indicator arrays that do not have the correct length and report an error (David Rader)

• Avoid triggering a libc assertion in contrib/hstore, due to use of memcpy() with equal source and destination pointers (Tomas Vondra)

• Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

The scripts in contrib/start-scripts/osx use infrastructure that's been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory contrib/start-scripts/macos containing scripts that use the newer launchd infrastructure.

• Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dunstan)

• Support linking to MinGW-built versions of libperl (Noah Misch)

This allows building PL/Perl with some common Perl distributions for Windows.

• Fix MSVC build to test whether 32-bit libperl needs -D_USE_32BIT_TIME_T (Noah Misch)

Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.

• On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)

This may allow collection of a core dump for some early-startup failures that did not produce a dump before.

• On Windows, avoid encoding-conversion-related crashes when emitting messages very early in postmaster startup (Takayuki Tsunakawa)

• Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)

• Add support for spinlocks on Motorola 88K (David Carlier)
• Update time zone data files to tzdata release 2018c for DST law changes in Brazil, Sao Tome and Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The US/Pacific-New zone has been removed (it was only an alias for America/Los_Angeles anyway).

E.64. Release 9.3.20

Release date: 2017-11-09

This release contains a variety of fixes from 9.3.19. For information about new features in the 9.3 major release, see Section E.84.

E.64.1. Migration to Version 9.3.20

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.18, see Section E.66.

E.64.2. Changes

• Fix crash due to rowtype mismatch in json{b}_populate_recordset() (Michael Paquier, Tom Lane)
  These functions used the result rowtype specified in the FROM ... AS clause without checking that it matched the actual rowtype of the supplied tuple value. If it didn't, that would usually result in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)

• Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)
  Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making $PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

  By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

• Properly reject attempts to convert infinite float values to type numeric (Tom Lane, KaiGai Kohei)
  Previously the behavior was platform-dependent.

• Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)

• Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes (Tom Lane)
  Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

• Correctly detect hashability of range data types (Tom Lane)
  The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range's subtype has hash support. This does not affect any of the built-in range types, since they're all hashable anyway.

• Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)
  If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.
• Prevent low-probability crash in processing of nested trigger firings (Tom Lane)

• Correctly restore the umask setting when file creation fails in COPY or lo_export() (Peter Eisentraut)

• Give a better error message for duplicate column names in ANALYZE (Nathan Bossart)

• Fix mis-parsing of the last line in a non-newline-terminated pg_hba.conf file (Tom Lane)

• Fix libpq to not require user's home directory to exist (Tom Lane)

In v10, failure to find the home directory while trying to read ~/.pgpass was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading ~/.pg_service.conf, though this was less obvious since that file is not sought unless a service name is specified.

• Fix libpq to guard against integer overflow in the row count of a PGresult (Michael Paquier)

• Fix ecpg's handling of out-of-scope cursor declarations with pointer or array variables (Michael Meskes)

• Make ecpglib's Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengqi, Michael Meskes)

• Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

This fixes various issues; the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no posixrules file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.

• Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

E.65. Release 9.3.19

Release date: 2017-08-31

This release contains a small number of fixes from 9.3.18. For information about new features in the 9.3 major release, see Section E.84.

E.65.1. Migration to Version 9.3.19

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.18, see Section E.66.

E.65.2. Changes

• Show foreign tables in information_schema.table_privileges view (Peter Eisentraut)

All other relevant information_schema views include foreign tables, but this one ignored them.

Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:
SET search_path TO information_schema;
CREATE OR REPLACE VIEW table_privileges AS
    SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
           CAST(grantee.rolname AS sql_identifier) AS grantee,
           CAST(current_database() AS sql_identifier) AS table_catalog,
           CAST(nc.nspname AS sql_identifier) AS table_schema,
           CAST(c.relname AS sql_identifier) AS table_name,
           CAST(c.prtype AS character_data) AS privilege_type,
           CAST(CASE WHEN
                       -- object owner always has grant options
                       pg_has_role(grantee.oid, c.relowner, 'USAGE')
                       OR c.grantable
                       THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
           CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE 'NO' END AS yes_or_no) AS with_hierarchy
    FROM (SELECT oid, relname, relnamespace, relkind, relowner,
                (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class
            ) AS c (oid, relname, relnamespace, relkind, relowner,
                    grantor, grantee, prtype, grantable),
           pg_namespace nc,
           pg_authid u_grantor,
           (SELECT oid, rolname FROM pg_authid
            UNION ALL
            SELECT 0::oid, 'PUBLIC'
           ) AS grantee (oid, rolname)
    WHERE c.relnamespace = nc.oid
        AND c.relkind IN ('r', 'v', 'f')
        AND c.grantee = grantee.oid
        AND c.grantor = u_grantor.oid
        AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE', 'DELETE', 'TRUNCATE', 'REFERENCES', 'TRIGGER')
        AND (pg_has_role(u_grantor.oid, 'USAGE')
             OR pg_has_role(grantee.oid, 'USAGE')
             OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

• Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)

  This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about “cannot drop active portal”.

• Remove assertion that could trigger during a fatal exit (Tom Lane)

• Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

  Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant val-
ues that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.

- Fix crash in pg_restore when using parallel mode and using a list file to select a subset of items to restore (Fabrízio de Royes Mello)
- Change ecpg's parser to allow RETURNING clauses without attached C variables (Michael Meskes)

This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

- Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that _USE_32BIT_TIME_T is no longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

**E.66. Release 9.3.18**

**Release date:** 2017-08-10

This release contains a variety of fixes from 9.3.17. For information about new features in the 9.3 major release, see Section E.84.

**E.66.1. Migration to Version 9.3.18**

A dump/restore is not required for those running 9.3.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.3.16, see Section E.68.

**E.66.2. Changes**

- Further restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Noah Misch)

The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have USAGE permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since information_schema.user_mapping_options does not show the options in such cases, pg_user_mappings should not either. (CVE-2017-7547)

By itself, this patch will only fix the behavior in newly initdb’d databases. If you wish to apply this change in an existing database, you will need to do the following:

1. Restart the postmaster after adding allow_system_table_mods = true to postgresql.conf. (In versions supporting ALTER SYSTEM, you can use that to make the configuration change, but you'll still need a restart.)

2. In each database of the cluster, run the following commands as superuser:

```sql
SET search_path = pg_catalog;
CREATE OR REPLACE VIEW pg_user_mappings AS
SELECT
```
Release Notes

U.oid AS umid,
S.oid AS srvid,
S.srvname AS srvname,
U.umuser AS umuser,
CASE WHEN U.umuser = 0 THEN 'public'
ELSE A.rolname
END AS usename,
CASE WHEN (U.umuser <> 0 AND A.rolname = current_user
AND (pg_has_role(S.srvowner, 'USAGE')
OR has_server_privilege(S.oid,
'USAGE')))
OR (U.umuser = 0 AND pg_has_role(S.srvowner, 'USAGE'))
OR (SELECT rolsuper FROM pg_authid WHERE
rolname = current_user)
THEN U.umoptions
ELSE NULL END AS umoptions
FROM pg_user_mapping U
LEFT JOIN pg_authid A ON (A.oid = U.umuser) JOIN
pg_foreign_server S ON (U.umserver = S.oid);

3. Do not forget to include the template0 and template1 databases, or the vulnerability will still exist in databases you create later. To fix template0, you'll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

```
ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;
```

and then after fixing template0, undo that with

```
ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;
```

In prior versions, instead use

```
UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';
```

4. Finally, remove the allow_system_table_mods configuration setting, and again restart the postmaster.

- Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)

libpq ignores empty password specifications, and does not transmit them to the server. So, if a user's password has been set to the empty string, it's impossible to log in with that password via psql or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, md5, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

- Fix concurrent locking of tuple update chains (Álvaro Herrera)

If several sessions concurrently lock a tuple update chain with nonconflicting lock modes using an old snapshot, and they all succeed, it was possible for some of them to nonetheless fail (and conclude

2627
there is no live tuple version) due to a race condition. This had consequences such as foreign-key
checks failing to see a tuple that definitely exists but is being updated concurrently.

• Fix potential data corruption when freezing a tuple whose XMAX is a multixact with exactly one
still-interesting member (Teodor Sigaev)

• On Windows, retry process creation if we fail to reserve the address range for our shared memory
in the new process (Tom Lane, Amit Kapila)

  This is expected to fix infrequent child-process-launch failures that are probably due to interference
from antivirus products.

• Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas
Munro, Tom Lane)

• Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael
Paquier)

• Prevent sending SSL session tickets to clients (Tom Lane)

  This fix prevents reconnection failures with ticket-aware client-side SSL code.

• Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)

• Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and im-
mediate restart (Tom Lane)

  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively
ignored.

• Ensure that the statistics collector's receive buffer size is at least 100KB (Tom Lane)

  This reduces the risk of dropped statistics data on older platforms whose default receive buffer size
is less than that.

• Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes
an XLOG_SWITCH WAL record (Andres Freund)

• Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)

• Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom
Lane)

• Fix cases where an INSERT or UPDATE assigns to more than one element of a column that is of
domain-over-array type (Tom Lane)

• Allow window functions to be used in sub-SELECTs that are within the arguments of an aggregate
function (Tom Lane)

• Move autogenerated array types out of the way during ALTER ... RENAME (Vik Fearing)

  Previously, we would rename a conflicting autogenerated array type out of the way during CREATE;
this fix extends that behavior to renaming operations.

• Ensure that ALTER USER ... SET accepts all the syntax variants that ALTER ROLE ... SET does (Peter Eisentraut)

• Properly update dependency info when changing a datatype I/O function's argument or return type
from opaque to the correct type (Heikki Linnakangas)

  CREATE TYPE updates I/O functions declared in this long-obsolete style, but it forgot to record a
dependency on the type, allowing a subsequent DROP TYPE to leave broken function definitions
behind.
• Reduce memory usage when ANALYZE processes a tsvector column (Heikki Linnakangas)

• Fix unnecessary precision loss and sloppy rounding when multiplying or dividing money values by integers or floats (Tom Lane)

• Tighten checks for whitespace in functions that parse identifiers, such as regprocedurein() (Tom Lane)

  Depending on the prevailing locale, these functions could misinterpret fragments of multibyte characters as whitespace.

• Use relevant #define symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)

  This avoids portability problems, typically manifesting as a “handshake” mismatch during library load, when working with recent Perl versions.

• In libpq, reset GSS/SASL and SSPI authentication state properly after a failed connection attempt (Michael Paquier)

  Failure to do this meant that when falling back from SSL to non-SSL connections, a GSS/SASL failure in the SSL attempt would always cause the non-SSL attempt to fail. SSPI did not fail, but it leaked memory.

• In psql, fix failure when COPY FROM STDIN is ended with a keyboard EOF signal and then another COPY FROM STDIN is attempted (Thomas Munro)

  This misbehavior was observed on BSD-derived platforms (including macOS), but not on most others.

• Fix pg_dump and pg_restore to emit REFRESH MATERIALIZED VIEW commands last (Tom Lane)

  This prevents errors during dump/restore when a materialized view refers to tables owned by a different user.

• Fix pg_dump with the --clean option to drop event triggers as expected (Tom Lane)

  It also now correctly assigns ownership of event triggers; before, they were restored as being owned by the superuser running the restore script.

• Fix pg_dump to not emit invalid SQL for an empty operator class (Daniel Gustafsson)

• Fix pg_dump output to stdout on Windows (Kuntal Ghosh)

  A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

• Fix pg_get_ruledef() to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)

  In some corner cases, pg_dump relies on pg_get_ruledef() to dump views, so that this error could result in dump/reload failures.

• Fix dumping of outer joins with empty constraints, such as the result of a NATURAL LEFT JOIN with no common columns (Tom Lane)

• Fix dumping of function expressions in the FROM clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)

• Fix pg_basebackup output to stdout on Windows (Haribabu Kommi)
A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix `pg_upgrade` to ensure that the ending WAL record does not have `wal_level = minimum` (Bruce Momjian)

  This condition could prevent upgraded standby servers from reconnecting.

- In `postgres_fdw`, re-establish connections to remote servers after `ALTER SERVER` or `ALTER USER MAPPING` commands (Kyotaro Horiguchi)

  This ensures that option changes affecting connection parameters will be applied promptly.

- In `postgres_fdw`, allow cancellation of remote transaction control commands (Robert Haas, Rafia Sabih)

  This change allows us to quickly escape a wait for an unresponsive remote server in many more cases than previously.

- Always use `-fPIC`, not `-fpic`, when building shared libraries with gcc (Tom Lane)

  This supports larger extension libraries on platforms where it makes a difference.

- Fix unescaped-braces issue in our build scripts for Microsoft MSVC, to avoid a warning or error from recent Perl versions (Andrew Dunstan)

- In MSVC builds, handle the case where the OpenSSL library is not within a VC subdirectory (Andrew Dunstan)

- In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)

  This fixes a former need to move things around in standard Windows installations of libxml2.

- In MSVC builds, recognize a Tcl library that is named `tcl86.lib` (Noah Misch)

---

**E.67. Release 9.3.17**

**Release date:** 2017-05-11

This release contains a variety of fixes from 9.3.16. For information about new features in the 9.3 major release, see Section E.84.

**E.67.1. Migration to Version 9.3.17**

A dump/restore is not required for those running 9.3.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.3.16, see Section E.68.

**E.67.2. Changes**

- Restrict visibility of `pg_user_mappings.umoptions`, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)

  The previous coding allowed the owner of a foreign server object, or anyone he has granted server `USAGE` permission to, to see the options for all user mappings associated with that server. This
might well include passwords for other users. Adjust the view definition to match the behavior of information_schema.user_mapping_options, namely that these options are visible to the user being mapped, or if the mapping is for PUBLIC and the current user is the server owner, or if the current user is a superuser. (CVE-2017-7486)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, follow the corrected procedure shown in the changelog entry for CVE-2017-7547, in Section E.66.

• Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

Some selectivity estimation functions in the planner will apply user-defined operators to values obtained from pg_statistic, such as most common values and histogram entries. This occurs before table permissions are checked, so a nefarious user could exploit the behavior to obtain these values for table columns he does not have permission to read. To fix, fall back to a default estimate if the operator's implementation function is not certified leak-proof and the calling user does not have permission to read the table column whose statistics are needed. At least one of these criteria is satisfied in most cases in practice. (CVE-2017-7484)

• Restore libpq's recognition of the PGREQUIRESSL environment variable (Daniel Gustafsson)

Processing of this environment variable was unintentionally dropped in PostgreSQL 9.3, but its documentation remained. This creates a security hazard, since users might be relying on the environment variable to force SSL-encrypted connections, but that would no longer be guaranteed. Restore handling of the variable, but give it lower priority than PGSSLMODE, to avoid breaking configurations that work correctly with post-9.3 code. (CVE-2017-7485)

• Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)

This could result in an unlogged index being set to an invalid state after a crash and restart. Such a problem would persist until the index was dropped and rebuilt.

• Fix incorrect reconstruction of pg_subtrans entries when a standby server replays a prepared but uncommitted two-phase transaction (Tom Lane)

In most cases this turned out to have no visible ill effects, but in corner cases it could result in circular references in pg_subtrans, potentially causing infinite loops in queries that examine rows modified by the two-phase transaction.

• Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL (Julien Rouhaud, Tom Lane)

Due to lack of a cache flush step between commands in an extension script file, non-utility queries might not see the effects of an immediately preceding catalog change, such as ALTER TABLE ... RENAME.

• Skip tablespace privilege checks when ALTER TABLE ... ALTER COLUMN TYPE rebuilds an existing index (Noah Misch)

The command failed if the calling user did not currently have CREATE privilege for the tablespace containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be rebuilt where it is.

• Fix ALTER TABLE ... VALIDATE CONSTRAINT to not recurse to child tables when the constraint is marked NO INHERIT (Amit Langote)

This fix prevents unwanted “constraint does not exist” failures when no matching constraint is present in the child tables.

• Fix VACUUM to account properly for pages that could not be scanned due to conflicting page pins (Andrew Gierth)
This tended to lead to underestimation of the number of tuples in the table. In the worst case of a small heavily-contended table, VACUUM could incorrectly report that the table contained no tuples, leading to very bad planning choices.

• Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests (Tom Lane, Thomas Munro)

• Fix cursor_to_xml() to produce valid output with tableforest = false (Thomas Munro, Peter Eisentraut)

Previously it failed to produce a wrapping <table> element.

• Improve performance of pg_timezone_names view (Tom Lane, David Rowley)

• Fix sloppy handling of corner-case errors from lseek() and close() (Tom Lane)

Neither of these system calls are likely to fail in typical situations, but if they did, fd.c could get quite confused.

• Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)

This could result in attempting to write to the event log when that isn't accessible, so that no logging happens at all.

• Fix ecpg to support COMMIT PREPARED and ROLLBACK PREPARED (Masahiko Sawada)

• Fix a double-free error when processing dollar-quoted string literals in ecpg (Michael Meskes)

• In pg_dump, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)

In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.

• Avoid emitting an invalid list file in pg_restore -l when SQL object names contain newlines (Tom Lane)

Replace newlines by spaces, which is sufficient to make the output valid for pg_restore -L's purposes.

• Fix pg_upgrade to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)

Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.

• Improve error handling in contrib/adminpack's pg_file_write() function (Noah Misch)

Notably, it failed to detect errors reported by fclose().

• In contrib/dblink, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

• Fix contrib/pg_trgm's extraction of trigrams from regular expressions (Tom Lane)

In some cases it would produce a broken data structure that could never match anything, leading to GIN or GiST index scans that use a trigram index not finding any matches to the regular expression.

• In contrib/postgres_fdw, transmit query cancellation requests to the remote server (Michael Paquier, Etsuro Fujita)
Previously, a local query cancellation request did not cause an already-sent remote query to terminate early. This is a back-patch of work originally done for 9.6.

- **Support OpenSSL 1.1.0** (Heikki Linnakangas, Andreas Karlsson, Tom Lane)
  
  This is a back-patch of work previously done in newer branches; it’s needed since many platforms are adopting newer OpenSSL versions.

- **Support Tcl 8.6 in MSVC builds** (Álvaro Herrera)

- **Sync our copy of the timezone library with IANA release tzcode2017b** (Tom Lane)
  
  This fixes a bug affecting some DST transitions in January 2038.

- **Update time zone data files to tzdata release 2017b for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.**

  The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the `pg_timezone_names` view nor used for output.

- **Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds** (David Rowley)

  The Microsoft MSVC build scripts neglected to install the `posixrules` file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the `US/Eastern` zone. If you don't want that, remove the `posixrules` file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.

---

**E.68. Release 9.3.16**

**Release date:** 2017-02-09

This release contains a variety of fixes from 9.3.15. For information about new features in the 9.3 major release, see Section E.84.

**E.68.1. Migration to Version 9.3.16**

A dump/restore is not required for those running 9.3.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.3.15, see Section E.69.

**E.68.2. Changes**

- Fix a race condition that could cause indexes built with `CREATE INDEX CONCURRENTLY` to be corrupt (Pavan Deolasee, Tom Lane)
If `CREATE INDEX CONCURRENTLY` was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the `CREATE INDEX` command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

- Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)

  Previously, this was skipped when `wal_level = minimal`, but actually it's necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

- If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)

- Ensure that hot standby feedback works correctly when it's enabled at standby server start (Ants Aasma, Craig Ringer)

- Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

- Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)

  This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

- Fix check for when an extension member object can be dropped (Tom Lane)

  Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

- Make sure `ALTER TABLE` preserves index tablespace assignments when rebuilding indexes (Tom Lane, Michael Paquier)

  Previously, non-default settings of `default_tablespace` could result in broken indexes.

- Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)

  This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.

- Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via `ALTER TABLE ... INHERIT` (Amit Langote)

  The OID column should be treated the same as regular user columns in this case, but it wasn't, leading to odd behavior in later inheritance changes.

- Report correct object identity during `ALTER TEXT SEARCH CONFIGURATION` (Artur Zakirov)

  The wrong catalog OID was reported to extensions such as logical decoding.

- Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

  When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

- Prevent multicolumn expansion of `foo.*` in an `UPDATE` source expression (Tom Lane)

  This led to “UPDATE target count mismatch --- internal error”. Now the syntax is understood as a whole-row variable, as it would be in other contexts.
• Ensure that column typmods are determined accurately for multi-row VALUES constructs (Tom Lane)

This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a varchar value) but later values don't share the same limit.

• Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (Uš '...') or Unicode identifier (Uš "...").

• Ensure that a purely negative text search query, such as !foo, matches empty tsvectors (Tom Dunstan)

Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

• Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

• Fix performance problems in ts_rewrite() (Tom Lane)

• Fix ts_rewrite()'s handling of nested NOT operators (Tom Lane)

• Fix array_fill() to handle empty arrays properly (Tom Lane)

• Fix one-byte buffer overrun in quote_literal_cstr() (Heikki Linnakangas)

The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.

• Prevent multiple calls of pg_start_backup() and pg_stop_backup() from running concurrently (Michael Paquier)

This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

• Avoid discarding interval-to-interval casts that aren't really no-ops (Tom Lane)

In some cases, a cast that should result in zeroing out low-order interval fields was mistakenly deemed to be a no-op and discarded. An example is that casting from INTERVAL MONTH to INTERVAL YEAR failed to clear the months field.

• Ensure that cached plans are invalidated by changes in foreign-table options (Amit Langote, Etsuro Fujita, Ashutosh Bapat)

• Fix pg_dump to dump user-defined casts and transforms that use built-in functions (Stephen Frost)

• Fix possible pg_basebackup failure on standby server when including WAL files (Amit Kapila, Robert Haas)

• Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)

This avoids failures if the objects are used after a Python garbage collection cycle has occurred.

• Fix PL/Tcl to support triggers on tables that have .tupno as a column name (Tom Lane)

This matches the (previously undocumented) behavior of PL/Tcl's spi_exec and spi_execp commands, namely that a magic .tupno column is inserted only if there isn't a real column named that.
• Allow DOS-style line endings in ~/.pgpass files, even on Unix (Vik Fearing)

  This change simplifies use of the same password file across Unix and Windows machines.

• Fix one-byte buffer overrun if ecpg is given a file name that ends with a dot (Takayuki Tsunakawa)

• Fix psql’s tab completion for \texttt{\texttt{ALTER DEFAULT PRIVILEGES}} (Gilles Darold, Stephen Frost)

• In psql, treat an empty or all-blank setting of the \texttt{PAGER} environment variable as meaning “no pager” (Tom Lane)

  Previously, such a setting caused output intended for the pager to vanish entirely.

• Improve \texttt{contrib/dblink}’s reporting of low-level libpq errors, such as out-of-memory (Joe Conway)

• Teach \texttt{contrib/dblink} to ignore irrelevant server options when it uses a \texttt{contrib/postgresql_fdw} foreign server as the source of connection options (Corey Huinker)

  Previously, if the foreign server object had options that were not also libpq connection options, an error occurred.

• On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)

• Sync our copy of the timezone library with IANA release tzcode2016j (Tom Lane)

  This fixes various issues, most notably that timezone data installation failed if the target directory didn’t support hard links.

• Update time zone data files to tzdata release 2016j for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarctica/Casey. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.

\section*{E.69. Release 9.3.15}

\textbf{Release date:} 2016-10-27

This release contains a variety of fixes from 9.3.14. For information about new features in the 9.3 major release, see Section E.84.

\section*{E.69.1. Migration to Version 9.3.15}

A dump/restore is not required for those running 9.3.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted free space maps.

Also, if you are upgrading from a version earlier than 9.3.9, see Section E.75.

\section*{E.69.2. Changes}

• Fix WAL-logging of truncation of relation free space maps and visibility maps (Pavan Deolasee, Heikki Linnakangas)

  It was possible for these files to not be correctly restored during crash recovery, or to be written incorrectly on a standby server. Bogus entries in a free space map could lead to attempts to access
pages that have been truncated away from the relation itself, typically producing errors like “could not read block XXX: read only 0 of 8192 bytes”. Checksum failures in the visibility map are also possible, if checksumming is enabled.

Procedures for determining whether there is a problem and repairing it if so are discussed at https://wiki.postgresql.org/wiki/Free_Space_Map_Problems.

- Fix SELECT FOR UPDATE/SHARE to correctly lock tuples that have been updated by a subsequently-aborted transaction (Álvaro Herrera)

  In 9.5 and later, the SELECT would sometimes fail to return such tuples at all. A failure has not been proven to occur in earlier releases, but might be possible with concurrent updates.

- Fix EvalPlanQual rechecks involving CTE scans (Tom Lane)

  The recheck would always see the CTE as returning no rows, typically leading to failure to update rows that were recently updated.

- Fix improper repetition of previous results from hashed aggregation in a subquery (Andrew Gierth)

  The test to see if we can reuse a previously-computed hash table of the aggregate state values neglected the possibility of an outer query reference appearing in an aggregate argument expression. A change in the value of such a reference should lead to recalculating the hash table, but did not.

- Fix EXPLAIN to emit valid XML when track_io_timing is on (Markus Winand)

  Previously the XML output-format option produced syntactically invalid tags such as `<I/O-Read-Time>`. That is now rendered as `<I-O-Read-Time>`.

- Suppress printing of zeroes for unmeasured times in EXPLAIN (Maksim Milyutin)

  Certain option combinations resulted in printing zero values for times that actually aren’t ever measured in that combination. Our general policy in EXPLAIN is not to print such fields at all, so do that consistently in all cases.

- Fix timeout length when VACUUM is waiting for exclusive table lock so that it can truncate the table (Simon Riggs)

  The timeout was meant to be 50 milliseconds, but it was actually only 50 microseconds, causing VACUUM to give up on truncation much more easily than intended. Set it to the intended value.

- Fix bugs in merging inherited CHECK constraints while creating or altering a table (Tom Lane, Amit Langote)

  Allow identical CHECK constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a NOT VALID constraint on the child. Likewise, prevent merging of a NO INHERIT child constraint with an inherited constraint.

- Remove artificial restrictions on the values accepted by numeric_in() and numeric_recv() (Tom Lane)

  We allow numeric values up to the limit of the storage format (more than 1e100000), so it seems fairly pointless that numeric_in() rejected scientific-notation exponents above 1000. Likewise, it was silly for numeric_recv() to reject more than 1000 digits in an input value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)

- Fix file descriptor leakage when truncating a temporary relation of more than 1GB (Andres Freund)

- Disallow starting a standalone backend with standby_mode turned on (Michael Paquier)
This can't do anything useful, since there will be no WAL receiver process to fetch more WAL data; and it could result in misbehavior in code that wasn't designed with this situation in mind.

• Don’t try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)
  This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)
  The reported problem involved leaking an error report during PQreset(), but there might be related cases.

• Make ecpg's --help and --version options work consistently with our other executables (Haribabu Kommi)

• In pg_dump, never dump range constructor functions (Tom Lane)
  This oversight led to pg_upgrade failures with extensions containing range types, due to duplicate creation of the constructor functions.

• In pg_xlogdump, retry opening new WAL segments when using --follow option (Magnus Hagander)
  This allows for a possible delay in the server's creation of the next segment.

• Fix pg_xlogdump to cope with a WAL file that begins with a continuation record spanning more than one page (Pavan Deolasee)

• Fix contrib/intarray/bench/bench.pl to print the results of the EXPLAIN it does when given the -e option (Daniel Gustafsson)

• Update Windows time zone mapping to recognize some time zone names added in recent Windows versions (Michael Paquier)

• Prevent failure of obsolete dynamic time zone abbreviations (Tom Lane)
  If a dynamic time zone abbreviation does not match any entry in the referenced time zone, treat it as equivalent to the time zone name. This avoids unexpected failures when IANA removes abbreviations from their time zone database, as they did in tzdata release 2016f and seem likely to do again in the future. The consequences were not limited to not recognizing the individual abbreviation; any mismatch caused the pg_timezone_abbrevs view to fail altogether.

• Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.70. Release 9.3.14

Release date: 2016-08-11
This release contains a variety of fixes from 9.3.13. For information about new features in the 9.3 major release, see Section E.84.

E.70.1. Migration to Version 9.3.14

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.9, see Section E.75.

E.70.2. Changes

• Fix possible mis-evaluation of nested CASE-WHEN expressions (Heikki Linnakangas, Michael Paquier, Tom Lane)

A CASE expression appearing within the test value subexpression of another CASE could become confused about whether its own test value was null or not. Also, inlining of a SQL function implementing the equality operator used by a CASE expression could result in passing the wrong test value to functions called within a CASE expression in the SQL function's body. If the test values were of different data types, a crash might result; moreover such situations could be abused to allow disclosure of portions of server memory. (CVE-2016-5423)

• Fix client programs' handling of special characters in database and role names (Noah Misch, Nathan Bossart, Michael Paquier)

Numerous places in vacuumdb and other client programs could become confused by database and role names containing double quotes or backslashes. Tighten up quoting rules to make that safe. Also, ensure that when a conninfo string is used as a database name parameter to these programs, it is correctly treated as such throughout.

Fix handling of paired double quotes in psql's \connect and \password commands to match the documentation.

Introduce a new -reuse-previous option in psql's \connect command to allow explicit control of whether to re-use connection parameters from a previous connection. (Without this, the choice is based on whether the database name looks like a conninfo string, as before.) This allows secure handling of database names containing special characters in pg_dumpall scripts.

pg_dumpall now refuses to deal with database and role names containing carriage returns or new-lines, as it seems impractical to quote those characters safely on Windows. In future we may reject such names on the server side, but that step has not been taken yet.

These are considered security fixes because crafted object names containing special characters could have been used to execute commands with superuser privileges the next time a superuser executes pg_dumpall or other routine maintenance operations. (CVE-2016-5424)

• Fix corner-case misbehaviors for IS NULL/IS NOT NULL applied to nested composite values (Andrew Gierth, Tom Lane)

The SQL standard specifies that IS NULL should return TRUE for a row of all null values (thus ROW(NULL,NULL) IS NULL yields TRUE), but this is not meant to apply recursively (thus ROW(NULL, ROW(NULL,NULL)) IS NULL yields FALSE). The core executor got this right, but certain planner optimizations treated the test as recursive (thus producing TRUE in both cases), and contrib/postgres_fdw could produce remote queries that misbehaved similarly.

• Make the inet and cidr data types properly reject IPv6 addresses with too many colon-separated fields (Tom Lane)

• Prevent crash in close_ps() (the point ## lseg operator) for NaN input coordinates (Tom Lane)
Release Notes

Make it return NULL instead of crashing.

- Avoid possible crash in \texttt{pg\_get\_expr()} when inconsistent values are passed to it (Michael Paquier, Thomas Munro)

- Fix several one-byte buffer over-reads in \texttt{to\_number()} (Peter Eisentraut)

  In several cases the \texttt{to\_number()} function would read one more character than it should from
  the input string. There is a small chance of a crash, if the input happens to be adjacent to the end
  of memory.

- Do not run the planner on the query contained in \texttt{CREATE MATERIALIZED VIEW} or \texttt{CREATE TABLE AS}
  when \texttt{WITH NO DATA} is specified (Michael Paquier, Tom Lane)

  This avoids some unnecessary failure conditions, for example if a stable function invoked by the
  materialized view depends on a table that doesn't exist yet.

- Avoid unsafe intermediate state during expensive paths through \texttt{heap\_update()} (Masahiko
  Sawada, Andres Freund)

  Previously, these cases locked the target tuple (by setting its XMAX) but did not WAL-log that
  action, thus risking data integrity problems if the page were spilled to disk and then a database crash
  occurred before the tuple update could be completed.

- Fix hint bit update during WAL replay of row locking operations (Andres Freund)

  The only known consequence of this problem is that row locks held by a prepared, but uncommitted,
  transaction might fail to be enforced after a crash and restart.

- Avoid unnecessary “could not serialize access” errors when acquiring \texttt{FOR KEY SHARE} row locks
  in serializable mode (Álvaro Herrera)

- Avoid crash in \texttt{postg}\texttt{res -C} when the specified variable has a null string value (Michael Paquier)

- Ensure that backends see up-to-date statistics for shared catalogs (Tom Lane)

  The statistics collector failed to update the statistics file for shared catalogs after a request from a
  regular backend. This problem was partially masked because the autovacuum launcher regularly
  makes requests that did cause such updates; however, it became obvious with autovacuum disabled.

- Avoid redundant writes of the statistics files when multiple backends request updates close together
  (Tom Lane, Tomas Vondra)

- Avoid consuming a transaction ID during \texttt{VACUUM} (Alexander Korotkov)

  Some cases in \texttt{VACUUM} unnecessarily caused an XID to be assigned to the current transaction.
  Normally this is negligible, but if one is up against the XID wraparound limit, consuming more
  XIDs during anti-wraparound vacuums is a very bad thing.

- Avoid canceling hot-standby queries during \texttt{VACUUM FREEZE} (Simon Riggs, Álvaro Herrera)

  \texttt{VACUUM FREEZE} on an otherwise-idle master server could result in unnecessary cancellations of
  queries on its standby servers.

- Prevent possible failure when vacuuming multixact IDs in an installation that has been \texttt{pg\_upgrade’d}
  from pre-9.3 (Andrew Gierth, Álvaro Herrera)

  The usual symptom of this bug is errors like “MultiXactId \texttt{NNN} has not been created yet -- apparent
  wraparound”.

- When a manual \texttt{ANALYZE} specifies a column list, don’t reset the table’s \texttt{changes\_since\_analyze}
  counter (Tom Lane)
Release Notes

If we’re only analyzing some columns, we should not prevent routine auto-analyze from happening for the other columns.

• Fix `ANALYZE`'s overestimation of `n_distinct` for a unique or nearly-unique column with many null entries (Tom Lane)

The nulls could get counted as though they were themselves distinct values, leading to serious planner misestimates in some types of queries.

• Prevent autovacuum from starting multiple workers for the same shared catalog (Álvaro Herrera)

Normally this isn’t much of a problem because the vacuum doesn’t take long anyway; but in the case of a severely bloated catalog, it could result in all but one worker uselessly waiting instead of doing useful work on other tables.

• Prevent infinite loop in GiST index build for geometric columns containing NaN component values (Tom Lane)

• Fix `contrib/btree_gin` to handle the smallest possible `bigint` value correctly (Peter Eisentraut)

• Teach `libpq` to correctly decode server version from future servers (Peter Eisentraut)

It’s planned to switch to two-part instead of three-part server version numbers for releases after 9.6. Make sure that `PQserverVersion()` returns the correct value for such cases.

• Fix `ecpg`'s code for `unsigned long long` array elements (Michael Meskes)

• In `pg_dump` with both `-c` and `-C` options, avoid emitting an unwanted `CREATE SCHEMA public` command (David Johnston, Tom Lane)

• Improve handling of SIGTERM/control-C in parallel `pg_dump` and `pg_restore` (Tom Lane)

Make sure that the worker processes will exit promptly, and also arrange to send query-cancel requests to the connected backends, in case they are doing something long-running such as a `CREATE INDEX`.

• Fix error reporting in parallel `pg_dump` and `pg_restore` (Tom Lane)

Previously, errors reported by `pg_dump` or `pg_restore` worker processes might never make it to the user's console, because the messages went through the master process, and there were various deadlock scenarios that would prevent the master process from passing on the messages. Instead, just print everything to `stderr`. In some cases this will result in duplicate messages (for instance, if all the workers report a server shutdown), but that seems better than no message.

• Ensure that parallel `pg_dump` or `pg_restore` on Windows will shut down properly after an error (Kyotaro Horiguchi)

Previously, it would report the error, but then just sit until manually stopped by the user.

• Make `pg_dump` behave better when built without zlib support (Kyotaro Horiguchi)

It didn’t work right for parallel dumps, and emitted some rather pointless warnings in other cases.

• Make `pg_basebackup` accept `-Z 0` as specifying no compression (Fujii Masao)

• Fix `makefiles`' rule for building AIX shared libraries to be safe for parallel make (Noah Misch)

• Fix TAP tests and MSVC scripts to work when build directory's path name contains spaces (Michael Paquier, Kyotaro Horiguchi)
• Be more predictable about reporting “statement timeout” versus “lock timeout” (Tom Lane)
  On heavily loaded machines, the regression tests sometimes failed due to reporting “lock timeout”
  even though the statement timeout should have occurred first.

• Make regression tests safe for Danish and Welsh locales (Jeff Janes, Tom Lane)
  Change some test data that triggered the unusual sorting rules of these locales.

• Update our copy of the timezone code to match IANA's tzcode release 2016c (Tom Lane)
  This is needed to cope with anticipated future changes in the time zone data files. It also fixes some
  corner-case bugs in coping with unusual time zones.

• Update time zone data files to tzdata release 2016f for DST law changes in Kemerovo and Novosibirsk,
  plus historical corrections for Azerbaijan, Belarus, and Morocco.

E.71. Release 9.3.13

Release date: 2016-05-12

This release contains a variety of fixes from 9.3.12. For information about new features in the 9.3 major release, see Section E.84.

E.71.1. Migration to Version 9.3.13

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.9, see Section E.75.

E.71.2. Changes

• Clear the OpenSSL error queue before OpenSSL calls, rather than assuming it's clear already; and
  make sure we leave it clear afterwards (Peter Geoghegan, Dave Vitek, Peter Eisentraut)
  This change prevents problems when there are multiple connections using OpenSSL within a single
  process and not all the code involved follows the same rules for when to clear the error queue.
  Failures have been reported specifically when a client application uses SSL connections in libpq
  concurrently with SSL connections using the PHP, Python, or Ruby wrappers for OpenSSL. It's
  possible for similar problems to arise within the server as well, if an extension module establishes
  an outgoing SSL connection.

• Fix “failed to build any n-way joins” planner error with a full join enclosed in the right-hand side
  of a left join (Tom Lane)

• Fix incorrect handling of equivalence-class tests in multilevel nestloop plans (Tom Lane)
  Given a three-or-more-way equivalence class of variables, such as \( X.X = Y.Y = Z.Z \), it was
  possible for the planner to omit some of the tests needed to enforce that all the variables are actually
  equal, leading to join rows being output that didn't satisfy the \( \text{WHERE} \) clauses. For various reasons,
  erroneous plans were seldom selected in practice, so that this bug has gone undetected for a long
  time.

• Fix possible misbehavior of \( \text{TH}, \text{th}, \text{and Y, YYY} \) format codes in \( \text{to_timestamp}() \) (Tom Lane)
  These could advance off the end of the input string, causing subsequent format codes to read
  garbage.

• Fix dumping of rules and views in which the \text{array} argument of a \text{value operator ANY (array)}
  construct is a sub-SELECT (Tom Lane)
• Make pg_regress use a startup timeout from the PGCTLTIMOUT environment variable, if that's set (Tom Lane)

This is for consistency with a behavior recently added to pg_ctl; it eases automated testing on slow machines.

• Fix pg_upgrade to correctly restore extension membership for operator families containing only one operator class (Tom Lane)

In such a case, the operator family was restored into the new database, but it was no longer marked as part of the extension. This had no immediate ill effects, but would cause later pg_dump runs to emit output that would cause (harmless) errors on restore.

• Fix pg_upgrade to not fail when new-cluster TOAST rules differ from old (Tom Lane)

pg_upgrade had special-case code to handle the situation where the new PostgreSQL version thinks that a table should have a TOAST table while the old version did not. That code was broken, so remove it, and instead do nothing in such cases; there seems no reason to believe that we can't get along fine without a TOAST table if that was okay according to the old version's rules.

• Back-port 9.4-era memory-barrier code changes into 9.2 and 9.3 (Tom Lane)

These changes were not originally needed in pre-9.4 branches, but we recently back-patched a fix that expected the barrier code to work properly. Only IA64 (when using icc), HPPA, and Alpha platforms are affected.

• Reduce the number of SysV semaphores used by a build configured with --disable-spinlocks (Tom Lane)

• Rename internal function strtoi() to strtoint() to avoid conflict with a NetBSD library function (Thomas Munro)

• Fix reporting of errors from bind() and listen() system calls on Windows (Tom Lane)

• Reduce verbosity of compiler output when building with Microsoft Visual Studio (Christian Ullrich)

• Fix putenv() to work properly with Visual Studio 2013 (Michael Paquier)

• Avoid possibly-unsafe use of Windows' FormatMessage() function (Christian Ullrich)

  Use the FORMAT_MESSAGE_IGNORE_INSERTS flag where appropriate. No live bug is known to exist here, but it seems like a good idea to be careful.

• Update time zone data files to tzdata release 2016d for DST law changes in Russia and Venezuela. There are new zone names Europe/Kirov and Asia/Tomsk to reflect the fact that these regions now have different time zone histories from adjacent regions.

E.72. Release 9.3.12

Release date: 2016-03-31

This release contains a variety of fixes from 9.3.11. For information about new features in the 9.3 major release, see Section E.84.

E.72.1. Migration to Version 9.3.12

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.9, see Section E.75.
E.72.2. Changes

- Fix incorrect handling of NULL index entries in indexed \texttt{ROW()} comparisons (Tom Lane)

An index search using a row comparison such as \texttt{ROW(a, b) > ROW('x', 'y')} would stop upon reaching a NULL entry in the \texttt{b} column, ignoring the fact that there might be non-NULL \texttt{b} values associated with later values of \texttt{a}.

- Avoid unlikely data-loss scenarios due to renaming files without adequate \texttt{fsync()} calls before and after (Michael Paquier, Tomas Vondra, Andres Freund)

- Correctly handle cases where \texttt{pg_subtrans} is close to XID wraparound during server startup (Jeff Janes)

- Fix corner-case crash due to trying to free \texttt{localeconv()} output strings more than once (Tom Lane)

- Fix parsing of affix files for \texttt{ispell} dictionaries (Tom Lane)

The code could go wrong if the affix file contained any characters whose byte length changes during case-folding, for example \texttt{ı} in Turkish UTF8 locales.

- Avoid use of \texttt{sscanf()} to parse \texttt{ispell} dictionary files (Artur Zakirov)

This dodges a portability problem on FreeBSD-derived platforms (including macOS).

- Avoid a crash on old Windows versions (before 7SP1/2008R2SP1) with an AVX2-capable CPU and a Postgres build done with Visual Studio 2013 (Christian Ullrich)

This is a workaround for a bug in Visual Studio 2013's runtime library, which Microsoft have stated they will not fix in that version.

- Fix \texttt{psql}'s tab completion logic to handle multibyte characters properly (Kyotaro Horiguchi, Robert Haas)

- Fix \texttt{psql}'s tab completion for \texttt{SECURITY LABEL} (Tom Lane)

Pressing \texttt{TAB} after \texttt{SECURITY LABEL} might cause a crash or offering of inappropriate keywords.

- Make \texttt{pg_ctl} accept a wait timeout from the \texttt{PGCTLTIMEOUT} environment variable, if none is specified on the command line (Noah Misch)

This eases testing of slower buildfarm members by allowing them to globally specify a longer-than-normal timeout for postmaster startup and shutdown.

- Fix incorrect test for Windows service status in \texttt{pg_ctl} (Manuel Mathar)

The previous set of minor releases attempted to fix \texttt{pg_ctl} to properly determine whether to send log messages to Window’s Event Log, but got the test backwards.

- Fix pgbench to correctly handle the combination of \texttt{-C} and \texttt{-M} prepared options (Tom Lane)

- In \texttt{pg_upgrade}, skip creating a deletion script when the new data directory is inside the old data directory (Bruce Momjian)

Blind application of the script in such cases would result in loss of the new data directory.

- In PL/Perl, properly translate empty Postgres arrays into empty Perl arrays (Alex Hunsaker)

- Make PL/Python cope with function names that aren't valid Python identifiers (Jim Nasby)
• Fix multiple mistakes in the statistics returned by contrib/pgstatu'spgstatindex() function (Tom Lane)

• Remove dependency on pseed in MSVC builds, since it's no longer provided by core Perl (Michael Paquier, Andrew Dunstan)

• Update time zone data files to tzdata release 2016c for DST law changes in Azerbaijan, Chile, Haiti, Palestine, and Russia (Altai, Astrakhan, Kirov, Sakhalin, Ulyanovsk regions), plus historical corrections for Lithuania, Moldova, and Russia (Kaliningrad, Samara, Volgograd).

E.73. Release 9.3.11

Release date: 2016-02-11

This release contains a variety of fixes from 9.3.10. For information about new features in the 9.3 major release, see Section E.84.

E.73.1. Migration to Version 9.3.11

A dump/restore is not required for those running 9.3.X. However, if you are upgrading from a version earlier than 9.3.9, see Section E.75.

E.73.2. Changes

• Fix infinite loops and buffer-overflow problems in regular expressions (Tom Lane)

  Very large character ranges in bracket expressions could cause infinite loops in some cases, and memory overwrites in other cases. (CVE-2016-0773)

• Perform an immediate shutdown if the postmaster.pid file is removed (Tom Lane)

  The postmaster now checks every minute or so that postmaster.pid is still there and still contains its own PID. If not, it performs an immediate shutdown, as though it had received SIGQUIT. The main motivation for this change is to ensure that failed buildfarm runs get cleaned up without manual intervention; but it also serves to limit the bad effects if a DBA forcibly removes postmaster.pid and then starts a new postmaster.

• In SERIALIZABLE transaction isolation mode, serialization anomalies could be missed due to race conditions during insertions (Kevin Grittner, Thomas Munro)

• Fix failure to emit appropriate WAL records when doing ALTER TABLE ... SET TABLESPACE for unlogged relations (Michael Paquier, Andres Freund)

  Even though the relation's data is unlogged, the move must be logged or the relation will be inaccessible after a standby is promoted to master.

• Fix possible misinitialization of unlogged relations at the end of crash recovery (Andres Freund, Michael Paquier)

• Ensure walsender slots are fully re-initialized when being re-used (Magnus Hagander)

• Fix ALTER COLUMN TYPE to reconstruct inherited check constraints properly (Tom Lane)

• Fix REASSIGN OWNED to change ownership of composite types properly (Álvaro Herrera)

• Fix REASSIGN OWNED and ALTER OWNER to correctly update granted-permissions lists when changing owners of data types, foreign data wrappers, or foreign servers (Bruce Momjian, Álvaro Herrera)
• Fix `REASSIGN OWNED` to ignore foreign user mappings, rather than fail (Álvaro Herrera)

• Fix possible crash after doing query rewrite for an updatable view (Stephen Frost)

• Fix planner’s handling of `LATERAL` references (Tom Lane)

  This fixes some corner cases that led to “failed to build any N-way joins” or “could not devise a query plan” planner failures.

• Add more defenses against bad planner cost estimates for GIN index scans when the index’s internal statistics are very out-of-date (Tom Lane)

• Make planner cope with hypothetical GIN indexes suggested by an index advisor plug-in (Julien Rouhaud)

• Speed up generation of unique table aliases in `EXPLAIN` and rule dumping, and ensure that generated aliases do not exceed `NAMEDATALEN` (Tom Lane)

• Fix dumping of whole-row Vars in `ROW()` and `VALUES()` lists (Tom Lane)

• Fix possible internal overflow in numeric division (Dean Rasheed)

• Fix enforcement of restrictions inside parentheses within regular expression lookahead constraints (Tom Lane)

  Lookahead constraints aren’t allowed to contain backrefs, and parentheses within them are always considered non-capturing, according to the manual. However, the code failed to handle these cases properly inside a parenthesized subexpression, and would give unexpected results.

• Conversion of regular expressions to indexscan bounds could produce incorrect bounds from regexps containing lookahead constraints (Tom Lane)

• Fix regular-expression compiler to handle loops of constraint arcs (Tom Lane)

  The code added for CVE-2007-4772 was both incomplete, in that it didn’t handle loops involving more than one state, and incorrect, in that it could cause assertion failures (though there seem to be no bad consequences of that in a non-assert build). Multi-state loops would cause the compiler to run until the query was canceled or it reached the too-many-states error condition.

• Improve memory-usage accounting in regular-expression compiler (Tom Lane)

  This causes the code to emit “regular expression is too complex” errors in some cases that previously used unreasonable amounts of time and memory.

• Improve performance of regular-expression compiler (Tom Lane)

• Make `%h` and `%r` escapes in `log_line_prefix` work for messages emitted due to `log_connections` (Tom Lane)

  Previously, `%h/%r` started to work just after a new session had emitted the “connection received” log message; now they work for that message too.

• On Windows, ensure the shared-memory mapping handle gets closed in child processes that don’t need it (Tom Lane, Amit Kapila)

  This oversight resulted in failure to recover from crashes whenever `logging_collector` is turned on.

• Fix possible failure to detect socket EOF in non-blocking mode on Windows (Tom Lane)

  It’s not entirely clear whether this problem can happen in pre-9.5 branches, but if it did, the symptom would be that a walsender process would wait indefinitely rather than noticing a loss of connection.
• Avoid leaking a token handle during SSPI authentication (Christian Ullrich)

• In psql, ensure that libreadline's idea of the screen size is updated when the terminal window size changes (Merlin Moncure)

Previously, libreadline did not notice if the window was resized during query output, leading to strange behavior during later input of multiline queries.

• Fix psql's \d\et command to interpret its pattern argument the same way as other \d commands with potentially schema-qualified patterns do (Reece Hart)

• Avoid possible crash in psql's \c command when previous connection was via Unix socket and command specifies a new hostname and same username (Tom Lane)

• In pg_ctl start -w, test child process status directly rather than relying on heuristics (Tom Lane, Michael Paquier)

Previously, pg_ctl relied on an assumption that the new postmaster would always create postmaster.pid within five seconds. But that can fail on heavily-loaded systems, causing pg_ctl to report incorrectly that the postmaster failed to start.

Except on Windows, this change also means that a pg_ctl start -w done immediately after another such command will now reliably fail, whereas previously it would report success if done within two seconds of the first command.

• In pg_ctl start -w, don't attempt to use a wildcard listen address to connect to the postmaster (Kondo Yuta)

On Windows, pg_ctl would fail to detect postmaster startup if listen_addresses is set to 0.0.0.0 or ::, because it would try to use that value verbatim as the address to connect to, which doesn't work. Instead assume that 127.0.0.1 or ::1, respectively, is the right thing to use.

• In pg_ctl on Windows, check service status to decide where to send output, rather than checking if standard output is a terminal (Michael Paquier)

• In pg_dump and pg_basebackup, adopt the GNU convention for handling tar-archive members exceeding 8GB (Tom Lane)

The POSIX standard for tar file format does not allow archive member files to exceed 8GB, but most modern implementations of tar support an extension that fixes that. Adopt this extension so that pg_dump with -Ft no longer fails on tables with more than 8GB of data, and so that pg_basebackup can handle files larger than 8GB. In addition, fix some portability issues that could cause failures for members between 4GB and 8GB on some platforms. Potentially these problems could cause unrecoverable data loss due to unreadable backup files.

• Fix assorted corner-case bugs in pg_dump's processing of extension member objects (Tom Lane)

• Make pg_dump mark a view's triggers as needing to be processed after its rule, to prevent possible failure during parallel pg_restore (Tom Lane)

• Ensure that relation option values are properly quoted in pg_dump (Kouhei Sutou, Tom Lane)

A reloption value that isn't a simple identifier or number could lead to dump/reload failures due to syntax errors in CREATE statements issued by pg_dump. This is not an issue with any reloption currently supported by core PostgreSQL, but extensions could allow reloptions that cause the problem.

• Avoid repeated password prompts during parallel pg_dump (Zeus Kronion)

• Fix pg_upgrade's file-copying code to handle errors properly on Windows (Bruce Momjian)
• Install guards in pgbench against corner-case overflow conditions during evaluation of script-specified division or modulo operators (Fabien Coelho, Michael Paquier)

• Fix failure to localize messages emitted by pg_receivexlog and pg_recvlogical (Joseph Kim)

• Avoid dump/reload problems when using both plpython2 and plpython3 (Tom Lane)

  In principle, both versions of PL/Python can be used in the same database, though not in the same session (because the two versions of libpython cannot safely be used concurrently). However, pg_restore and pg_upgrade both do things that can fall foul of the same-session restriction. Work around that by changing the timing of the check.

• Fix PL/Python regression tests to pass with Python 3.5 (Peter Eisentraut)

• Fix premature clearing of libpq’s input buffer when socket EOF is seen (Tom Lane)

  This mistake caused libpq to sometimes not report the backend’s final error message before reporting “server closed the connection unexpectedly”.

• Prevent certain PL/Java parameters from being set by non-superusers (Noah Misch)

  This change mitigates a PL/Java security bug (CVE-2016-0766), which was fixed in PL/Java by marking these parameters as superuser-only. To fix the security hazard for sites that update PostgreSQL more frequently than PL/Java, make the core code aware of them also.

• Improve libpq’s handling of out-of-memory situations (Michael Paquier, Amit Kapila, Heikki Linnakangas)

• Fix order of arguments in ecpg-generated typedef statements (Michael Meskes)

• Use %g not %f format in ecpg’s PGTYPESnumeric_from_double() (Tom Lane)

• Fix ecpg-supplied header files to not contain comments continued from a preprocessor directive line onto the next line (Michael Meskes)

  Such a comment is rejected by ecpg. It’s not yet clear whether ecpg itself should be changed.

• Fix hstore_to_json_loose()’s test for whether an hstore value can be converted to a JSON number (Tom Lane)

  Previously this function could be fooled by non-alphanumeric trailing characters, leading to emitting syntactically-invalid JSON.

• Ensure that contrib/pgcrypto’s crypt() function can be interrupted by query cancel (Andreas Karlsson)

• Accept flex versions later than 2.5.x (Tom Lane, Michael Paquier)

  Now that flex 2.6.0 has been released, the version checks in our build scripts needed to be adjusted.

• Improve reproducibility of build output by ensuring filenames are given to the linker in a fixed order (Christoph Berg)

  This avoids possible bitwise differences in the produced executable files from one build to the next.

• Install our missing script where PGXS builds can find it (Jim Nasby)

  This allows sane behavior in a PGXS build done on a machine where build tools such as bison are missing.

• Ensure that dynloader.h is included in the installed header files in MSVC builds (Bruce Momjian, Michael Paquier)
• Add variant regression test expected-output file to match behavior of current libxml2 (Tom Lane)

The fix for libxml2's CVE-2015-7499 causes it not to output error context reports in some cases where it used to do so. This seems to be a bug, but we'll probably have to live with it for some time, so work around it.

• Update time zone data files to tzdata release 2016a for DST law changes in Cayman Islands, Metlakatala, and Trans-Baikal Territory (Zabaykalsky Krai), plus historical corrections for Pakistan.

E.74. Release 9.3.10

Release date: 2015-10-08

This release contains a variety of fixes from 9.3.9. For information about new features in the 9.3 major release, see Section E.84.

E.74.1. Migration to Version 9.3.10

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.9, see Section E.75.

E.74.2. Changes

• Guard against stack overflows in json parsing (Oskari Saarenmaa)

If an application constructs PostgreSQL json or jsonb values from arbitrary user input, the application's users can reliably crash the PostgreSQL server, causing momentary denial of service. (CVE-2015-5289)

• Fix contrib/pgcrypto to detect and report too-short crypt() salts (Josh Kupershmidt)

Certain invalid salt arguments crashed the server or disclosed a few bytes of server memory. We have not ruled out the viability of attacks that arrange for presence of confidential information in the disclosed bytes, but they seem unlikely. (CVE-2015-5288)

• Fix subtransaction cleanup after a portal (cursor) belonging to an outer subtransaction fails (Tom Lane, Michael Paquier)

A function executed in an outer-subtransaction cursor could cause an assertion failure or crash by referencing a relation created within an inner subtransaction.

• Ensure all relations referred to by an updatable view are properly locked during an update statement (Dean Rasheed)

• Fix insertion of relations into the relation cache “init file” (Tom Lane)

An oversight in a patch in the most recent minor releases caused pg_trigger_tgrelid_tgname_index to be omitted from the init file. Subsequent sessions detected this, then deemed the init file to be broken and silently ignored it, resulting in a significant degradation in session startup time. In addition to fixing the bug, install some guards so that any similar future mistake will be more obvious.

• Avoid O(N^2) behavior when inserting many tuples into a SPI query result (Neil Conway)

• Improve LISTEN startup time when there are many unread notifications (Matt Newell)

• Fix performance problem when a session alters large numbers of foreign key constraints (Jan Wieck, Tom Lane)
This was seen primarily when restoring pg_dump output for databases with many thousands of tables.

- Disable SSL renegotiation by default (Michael Paquier, Andres Freund)

  While use of SSL renegotiation is a good idea in theory, we have seen too many bugs in practice, both in the underlying OpenSSL library and in our usage of it. Renegotiation will be removed entirely in 9.5 and later. In the older branches, just change the default value of `ssl_renegotiation_limit` to zero (disabled).

- Lower the minimum values of the *_freeze_max_age parameters (Andres Freund)

  This is mainly to make tests of related behavior less time-consuming, but it may also be of value for installations with limited disk space.

- Limit the maximum value of `wal_buffers` to 2GB to avoid server crashes (Josh Berkus)

- Avoid logging complaints when a parameter that can only be set at server start appears multiple times in `postgresql.conf`, and fix counting of line numbers after an `include_dir` directive (Tom Lane)

- Fix rare internal overflow in multiplication of `numeric` values (Dean Rasheed)

- Guard against hard-to-reach stack overflows involving record types, range types, `json`, `jsonb`, `tsquery`, `ltxtquery` and `query_int` (Noah Misch)

- Fix handling of `DOW` and `DOY` in datetime input (Greg Stark)

  These tokens aren't meant to be used in datetime values, but previously they resulted in opaque internal error messages rather than "invalid input syntax".

- Add more query-cancel checks to regular expression matching (Tom Lane)

- Add recursion depth protections to regular expression, `SIMILAR TO`, and `LIKE` matching (Tom Lane)

  Suitable search patterns and a low stack depth limit could lead to stack-overflow crashes.

- Fix potential infinite loop in regular expression execution (Tom Lane)

  A search pattern that can apparently match a zero-length string, but actually doesn't match because of a back reference, could lead to an infinite loop.

- In regular expression execution, correctly record match data for capturing parentheses within a quantifier even when the match is zero-length (Tom Lane)

- Fix low-memory failures in regular expression compilation (Andreas Seltenreich)

- Fix low-probability memory leak during regular expression execution (Tom Lane)

- Fix rare low-memory failure in lock cleanup during transaction abort (Tom Lane)

- Fix “unexpected out-of-memory situation during sort” errors when using tuplestores with small `work_mem` settings (Tom Lane)

- Fix very-low-probability stack overrun in `qsort` (Tom Lane)

- Fix “invalid memory alloc request size” failure in hash joins with large `work_mem` settings (Tomas Vondra, Tom Lane)

- Fix assorted planner bugs (Tom Lane)
These mistakes could lead to incorrect query plans that would give wrong answers, or to assertion failures in assert-enabled builds, or to odd planner errors such as “could not devise a query plan for the given query”, “could not find pathkey item to sort”, “plan should not reference subplan's variable”, or “failed to assign all NestLoopParams to plan nodes”. Thanks are due to Andreas Seltenreich and Piotr Stefaniak for fuzz testing that exposed these problems.

- Improve planner's performance for UPDATE/DELETE on large inheritance sets (Tom Lane, Dean Rasheed)

- Ensure standby promotion trigger files are removed at postmaster startup (Michael Paquier, Fujii Masao)

  This prevents unwanted promotion from occurring if these files appear in a database backup that is used to initialize a new standby server.

- During postmaster shutdown, ensure that per-socket lock files are removed and listen sockets are closed before we remove the postmaster.pid file (Tom Lane)

  This avoids race-condition failures if an external script attempts to start a new postmaster as soon as pg_ctl stop returns.

- Fix postmaster's handling of a startup-process crash during crash recovery (Tom Lane)

  If, during a crash recovery cycle, the startup process crashes without having restored database consistency, we'd try to launch a new startup process, which typically would just crash again, leading to an infinite loop.

- Make emergency autovacuuming for multixact wraparound more robust (Andres Freund)

- Do not print a WARNING when an autovacuum worker is already gone when we attempt to signal it, and reduce log verbosity for such signals (Tom Lane)

- Prevent autovacuum launcher from sleeping unduly long if the server clock is moved backwards a large amount (Álvaro Herrera)

- Ensure that cleanup of a GIN index's pending-insertions list is interruptable by cancel requests (Jeff Janes)

- Allow all-zeroes pages in GIN indexes to be reused (Heikki Linnakangas)

  Such a page might be left behind after a crash.

- Fix handling of all-zeroes pages in SP-GiST indexes (Heikki Linnakangas)

  VACUUM attempted to recycle such pages, but did so in a way that wasn't crash-safe.

- Fix off-by-one error that led to otherwise-harmless warnings about "apparent wraparound" in subtrans/multixact truncation (Thomas Munro)

- Fix misreporting of CONTINUE and MOVE statement types in PL/pgSQL's error context messages (Pavel Stehule, Tom Lane)

- Fix PL/Perl to handle non-ASCII error message texts correctly (Alex Hunsaker)

- Fix PL/Python crash when returning the string representation of a record result (Tom Lane)

- Fix some places in PL/Tcl that neglected to check for failure of malloc() calls (Michael Paquier, Álvaro Herrera)

- In contrib/isn, fix output of ISBN-13 numbers that begin with 979 (Fabien Coelho)
EANs beginning with 979 (but not 9790) are considered ISBNs, but they must be printed in the new 13-digit format, not the 10-digit format.

- Improve contrib/postgres_fdw's handling of collation-related decisions (Tom Lane)
  The main user-visible effect is expected to be that comparisons involving varchar columns will be sent to the remote server for execution in more cases than before.

- Improve libpq's handling of out-of-memory conditions (Michael Paquier, Heikki Linnakangas)

- Fix memory leaks and missing out-of-memory checks in ecpg (Michael Paquier)

- Fix psql's code for locale-aware formatting of numeric output (Tom Lane)
  The formatting code invoked by \pset numericlocale on did the wrong thing for some uncommon cases such as numbers with an exponent but no decimal point. It could also mangle already-localized output from the money data type.

- Prevent crash in psql's \c command when there is no current connection (Noah Misch)

- Make pg_dump handle inherited NOT VALID check constraints correctly (Tom Lane)

- Fix selection of default zlib compression level in pg_dump's directory output format (Andrew Dunstan)

- Ensure that temporary files created during a pg_dump run with tar-format output are not world-readable (Michael Paquier)

- Fix pg_dump and pg_upgrade to support cases where the postgres or template1 database is in a non-default tablespace (Marti Raudsepp, Bruce Momjian)

- Fix pg_dump to handle object privileges sanely when dumping from a server too old to have a particular privilege type (Tom Lane)
  When dumping data types from pre-9.2 servers, and when dumping functions or procedural languages from pre-7.3 servers, pg_dump would produce GRANT/REVOKE commands that revoked the owner's grantable privileges and instead granted all privileges to PUBLIC. Since the privileges involved are just USAGE and EXECUTE, this isn't a security problem, but it's certainly a surprising representation of the older systems' behavior. Fix it to leave the default privilege state alone in these cases.

- Fix pg_dump to dump shell types (Tom Lane)
  Shell types (that is, not-yet-fully-defined types) aren't useful for much, but nonetheless pg_dump should dump them.

- Fix assorted minor memory leaks in pg_dump and other client-side programs (Michael Paquier)

- Fix spinlock assembly code for PPC hardware to be compatible with AIX's native assembler (Tom Lane)
  Building with gcc didn't work if gcc had been configured to use the native assembler, which is becoming more common.

- On AIX, test the -qlonglong compiler option rather than just assuming it's safe to use (Noah Misch)

- On AIX, use -Wl,-brtllib link option to allow symbols to be resolved at runtime (Noah Misch)
  Perl relies on this ability in 5.8.0 and later.

- Avoid use of inline functions when compiling with 32-bit xlc, due to compiler bugs (Noah Misch)
• Use librt for sched_yield() when necessary, which it is on some Solaris versions (Oskari Saarenmaa)

• Fix Windows install.bat script to handle target directory names that contain spaces (Heikki Linnakangas)

• Make the numeric form of the PostgreSQL version number (e.g., 90405) readily available to extension Makefiles, as a variable named VERSION_NUM (Michael Paquier)

• Update time zone data files to tzdata release 2015g for DST law changes in Cayman Islands, Fiji, Moldova, Morocco, Norfolk Island, North Korea, Turkey, and Uruguay. There is a new zone name America/Fort_Nelson for the Canadian Northern Rockies.

E.75. Release 9.3.9

Release date: 2015-06-12

This release contains a small number of fixes from 9.3.8. For information about new features in the 9.3 major release, see Section E.84.

E.75.1. Migration to Version 9.3.9

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading an installation that was previously upgraded using a pg_upgrade version between 9.3.0 and 9.3.4 inclusive, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.3.7, see Section E.77.

E.75.2. Changes

• Fix possible failure to recover from an inconsistent database state (Robert Haas)

Recent PostgreSQL releases introduced mechanisms to protect against multixact wraparound, but some of that code did not account for the possibility that it would need to run during crash recovery, when the database may not be in a consistent state. This could result in failure to restart after a crash, or failure to start up a secondary server. The lingering effects of a previously-fixed bug in pg_upgrade could also cause such a failure, in installations that had used pg_upgrade versions between 9.3.0 and 9.3.4.

The pg_upgrade bug in question was that it would set oldestMultiXid to 1 in pg_control even if the true value should be higher. With the fixes introduced in this release, such a situation will result in immediate emergency autovacuuming until a correct oldestMultiXid value can be determined. If that would pose a hardship, users can avoid it by doing manual vacuuming before upgrading to this release. In detail:

1. Check whether pg_controldata reports “Latest checkpoint's oldestMultiXid” to be 1. If not, there's nothing to do.

2. Look in PGDATA/pg_multixact/offsets to see if there's a file named 0000. If there is, there's nothing to do.

3. Otherwise, for each table that has pg_class.relminmxid equal to 1, VACUUM that table with both vacuum_multixact_freeze_min_age and vacuum_multixact_freeze_table_age set to zero. (You can use the vacuum cost delay parameters described in Section 19.4.4 to reduce the performance consequences for concurrent sessions.) You must use PostgreSQL 9.3.5 or later to perform this step.

• Fix rare failure to invalidate relation cache init file (Tom Lane)
With just the wrong timing of concurrent activity, a `VACUUM FULL` on a system catalog might fail to update the "init file" that's used to avoid cache-loading work for new sessions. This would result in later sessions being unable to access that catalog at all. This is a very ancient bug, but it's so hard to trigger that no reproducible case had been seen until recently.

- Avoid deadlock between incoming sessions and `CREATE/DROP DATABASE` (Tom Lane)

A new session starting in a database that is the target of a `DROP DATABASE` command, or is the template for a `CREATE DATABASE` command, could cause the command to wait for five seconds and then fail, even if the new session would have exited before that.

- Improve planner's cost estimates for semi-joins and anti-joins with inner index-scans (Tom Lane, Tomas Vondra)

This type of plan is quite cheap when all the join clauses are used as index scan conditions, even if the inner scan would nominally fetch many rows, because the executor will stop after obtaining one row. The planner only partially accounted for that effect, and would therefore overestimate the cost, leading it to possibly choose some other much less efficient plan type.

**E.76. Release 9.3.8**

**Release date:** 2015-06-04

This release contains a small number of fixes from 9.3.7. For information about new features in the 9.3 major release, see Section E.84.

**E.76.1. Migration to Version 9.3.8**

A dump/restore is not required for those running 9.3.X.

However, if you are upgrading from a version earlier than 9.3.7, see Section E.77.

**E.76.2. Changes**

- Avoid failures while `fsync`'ing data directory during crash restart (Abhijit Menon-Sen, Tom Lane)

In the previous minor releases we added a patch to `fsync` everything in the data directory after a crash. Unfortunately its response to any error condition was to fail, thereby preventing the server from starting up, even when the problem was quite harmless. An example is that an unwritable file in the data directory would prevent restart on some platforms; but it is common to make SSL certificate files unwritable by the server. Revise this behavior so that permissions failures are ignored altogether, and other types of failures are logged but do not prevent continuing.

Also apply the same rules in `initdb --sync-only`. This case is less critical but it should act similarly.

- Fix `pg_get_functiondef()` to show functions' `LEAKPROOF` property, if set (Jeevan Chalke)

- Remove configure's check prohibiting linking to a threaded libpython on OpenBSD (Tom Lane)

The failure this restriction was meant to prevent seems to not be a problem anymore on current OpenBSD versions.

- Allow `libpq` to use TLS protocol versions beyond v1 (Noah Misch)

For a long time, `libpq` was coded so that the only SSL protocol it would allow was TLS v1. Now that newer TLS versions are becoming popular, allow it to negotiate the highest commonly-supported TLS version with the server. (PostgreSQL servers were already capable of such negotiation, so no change is needed on the server side.) This is a back-patch of a change already released in 9.4.0.
E.77. Release 9.3.7

Release date: 2015-05-22

This release contains a variety of fixes from 9.3.6. For information about new features in the 9.3 major release, see Section E.84.

E.77.1. Migration to Version 9.3.7

A dump/restore is not required for those running 9.3.X.

However, if you use contrib/citext's regexp_matches() functions, see the changelog entry below about that.

Also, if you are upgrading from a version earlier than 9.3.6, see Section E.78.

E.77.2. Changes

- Avoid possible crash when client disconnects just before the authentication timeout expires (Benkocs Norbert Attila)

  If the timeout interrupt fired partway through the session shutdown sequence, SSL-related state would be freed twice, typically causing a crash and hence denial of service to other sessions. Experimentation shows that an unauthenticated remote attacker could trigger the bug somewhat consistently, hence treat as security issue. (CVE-2015-3165)

- Improve detection of system-call failures (Noah Misch)

  Our replacement implementation of snprintf() failed to check for errors reported by the underlying system library calls; the main case that might be missed is out-of-memory situations. In the worst case this might lead to information exposure, due to our code assuming that a buffer had been overwritten when it hadn't been. Also, there were a few places in which security-relevant calls of other system library functions did not check for failure.

  It remains possible that some calls of the *printf() family of functions are vulnerable to information disclosure if an out-of-memory error occurs at just the wrong time. We judge the risk to not be large, but will continue analysis in this area. (CVE-2015-3166)

- In contrib/pgcrypto, uniformly report decryption failures as “Wrong key or corrupt data” (Noah Misch)

  Previously, some cases of decryption with an incorrect key could report other error message texts. It has been shown that such variance in error reports can aid attackers in recovering keys from other systems. While it's unknown whether pgcrypto's specific behaviors are likewise exploitable, it seems better to avoid the risk by using a one-size-fits-all message. (CVE-2015-3167)

- Protect against wraparound of multixact member IDs (Álvaro Herrera, Robert Haas, Thomas Munro)

  Under certain usage patterns, the existing defenses against this might be insufficient, allowing pg_multixact/members files to be removed too early, resulting in data loss. The fix for this includes modifying the server to fail transactions that would result in overwriting old multixact member ID data, and improving autovacuum to ensure it will act proactively to prevent multixact member ID wraparound, as it does for transaction ID wraparound.

- Fix incorrect declaration of contrib/citext's regexp_matches() functions (Tom Lane)

  These functions should return setof text[], like the core functions they are wrappers for; but they were incorrectly declared as returning just text[]. This mistake had two results: first, if there
was no match you got a scalar null result, whereas what you should get is an empty set (zero rows). Second, the g flag was effectively ignored, since you would get only one result array even if there were multiple matches.

While the latter behavior is clearly a bug, there might be applications depending on the former behavior; therefore the function declarations will not be changed by default until PostgreSQL 9.5. In pre-9.5 branches, the old behavior exists in version 1.0 of the citext extension, while we have provided corrected declarations in version 1.1 (which is not installed by default). To adopt the fix in pre-9.5 branches, execute ALTER EXTENSION citext UPDATE TO '1.1' in each database in which citext is installed. (You can also “update” back to 1.0 if you need to undo that.) Be aware that either update direction will require dropping and recreating any views or rules that use citext's regexp_matches() functions.

• Fix incorrect checking of deferred exclusion constraints after a HOT update (Tom Lane)

If a new row that potentially violates a deferred exclusion constraint is HOT-updated (that is, no indexed columns change and the row can be stored back onto the same table page) later in the same transaction, the exclusion constraint would be reported as violated when the check finally occurred, even if the row(s) the new row originally conflicted with had been deleted.

• Fix planning of star-schema-style queries (Tom Lane)

Sometimes, efficient scanning of a large table requires that index parameters be provided from more than one other table (commonly, dimension tables whose keys are needed to index a large fact table). The planner should be able to find such plans, but an overly restrictive search heuristic prevented it.

• Prevent improper reordering of antijoins (NOT EXISTS joins) versus other outer joins (Tom Lane)

This oversight in the planner has been observed to cause “could not find RelOptInfo for given relids” errors, but it seems possible that sometimes an incorrect query plan might get past that consistency check and result in silently-wrong query output.

• Fix incorrect matching of subexpressions in outer-join plan nodes (Tom Lane)

Previously, if textually identical non-strict subexpressions were used both above and below an outer join, the planner might try to re-use the value computed below the join, which would be incorrect because the executor would force the value to NULL in case of an unmatched outer row.

• Fix GEQO planner to cope with failure of its join order heuristic (Tom Lane)

This oversight has been seen to lead to “failed to join all relations together” errors in queries involving LATERAL, and that might happen in other cases as well.

• Fix possible deadlock at startup when max_prepared_transactions is too small (Heikki Linnakangas)

• Don’t archive useless preallocated WAL files after a timeline switch (Heikki Linnakangas)

• Recursively fsync() the data directory after a crash (Abhijit Menon-Sen, Robert Haas)

This ensures consistency if another crash occurs shortly later. (The second crash would have to be a system-level crash, not just a database crash, for there to be a problem.)

• Fix autovacuum launcher's possible failure to shut down, if an error occurs after it receives SIGTERM (Álvaro Herrera)

• Cope with unexpected signals in LockBufferForCleanup() (Andres Freund)

This oversight could result in spurious errors about “multiple backends attempting to wait for pin-count 1”.

2656
• Fix crash when doing COPY IN to a table with check constraints that contain whole-row references (Tom Lane)

The known failure case only crashes in 9.4 and up, but there is very similar code in 9.3 and 9.2, so back-patch those branches as well.

• Avoid waiting for WAL flush or synchronous replication during commit of a transaction that was read-only so far as the user is concerned (Andres Freund)

Previously, a delay could occur at commit in transactions that had written WAL due to HOT page pruning, leading to undesirable effects such as sessions getting stuck at startup if all synchronous replicas are down. Sessions have also been observed to get stuck in catchup interrupt processing when using synchronous replication; this will fix that problem as well.

• Fix crash when manipulating hash indexes on temporary tables (Heikki Linnakangas)

• Fix possible failure during hash index bucket split, if other processes are modifying the index concurrently (Tom Lane)

• Check for interrupts while analyzing index expressions (Jeff Janes)

  ANALYZE executes index expressions many times; if there are slow functions in such an expression, it’s desirable to be able to cancel the ANALYZE before that loop finishes.

• Ensure tableoid of a foreign table is reported correctly when a READ COMMITTED recheck occurs after locking rows in SELECT FOR UPDATE, UPDATE, or DELETE (Etsuro Fujita)

• Add the name of the target server to object description strings for foreign-server user mappings (Álvaro Herrera)

• Include the schema name in object identity strings for conversions (Álvaro Herrera)

• Recommend setting include_realm to 1 when using Kerberos/GSSAPI/SSPI authentication (Stephen Frost)

  Without this, identically-named users from different realms cannot be distinguished. For the moment this is only a documentation change, but it will become the default setting in PostgreSQL 9.5.

• Remove code for matching IPv4 pg_hba.conf entries to IPv4-in-IPv6 addresses (Tom Lane)

  This hack was added in 2003 in response to a report that some Linux kernels of the time would report IPv4 connections as having IPv4-in-IPv6 addresses. However, the logic was accidentally broken in 9.0. The lack of any field complaints since then shows that it’s not needed anymore. Now we have reports that the broken code causes crashes on some systems, so let’s just remove it rather than fix it. (Had we chosen to fix it, that would make for a subtle and potentially security-sensitive change in the effective meaning of IPv4 pg_hba.conf entries, which does not seem like a good thing to do in minor releases.)

• Report WAL flush, not insert, position in IDENTIFY_SYSTEM replication command (Heikki Linnakangas)

  This avoids a possible startup failure in pg_receivexlog.

• While shutting down service on Windows, periodically send status updates to the Service Control Manager to prevent it from killing the service too soon; and ensure that pg_ctl will wait for shutdown (Krystian Bigaj)

• Reduce risk of network deadlock when using libpq’s non-blocking mode (Heikki Linnakangas)

  When sending large volumes of data, it’s important to drain the input buffer every so often, in case the server has sent enough response data to cause it to block on output. (A typical scenario is that
the server is sending a stream of NOTICE messages during COPY FROM STDIN.) This worked properly in the normal blocking mode, but not so much in non-blocking mode. We've modified libpq to opportunistically drain input when it can, but a full defense against this problem requires application cooperation: the application should watch for socket read-ready as well as write-ready conditions, and be sure to call PQconsumeInput() upon read-ready.

- In libpq, fix misparsing of empty values in URI connection strings (Thomas Fanghaenel)
- Fix array handling in ecpg (Michael Meskes)
- Fix psql to sanely handle URIs and conninfo strings as the first parameter to \connect (David Fetter, Andrew Dunstan, Álvaro Herrera)

This syntax has been accepted (but undocumented) for a long time, but previously some parameters might be taken from the old connection instead of the given string, which was agreed to be undesirable.

- Suppress incorrect complaints from psql on some platforms that it failed to write ~/.psql_history at exit (Tom Lane)

This misbehavior was caused by a workaround for a bug in very old (pre-2006) versions of libedit. We fixed it by removing the workaround, which will cause a similar failure to appear for anyone still using such versions of libedit. Recommendation: upgrade that library, or use libreadline.

- Fix pg_dump's rule for deciding which casts are system-provided casts that should not be dumped (Tom Lane)
- In pg_dump, fix failure to honor -Z compression level option together with -Fd (Michael Paquier)
- Make pg_dump consider foreign key relationships between extension configuration tables while choosing dump order (Gilles Darold, Michael Paquier, Stephen Frost)

This oversight could result in producing dumps that fail to reload because foreign key constraints are transiently violated.

- Avoid possible pg_dump failure when concurrent sessions are creating and dropping temporary functions (Tom Lane)
- Fix dumping of views that are just VALUES(...) but have column aliases (Tom Lane)
- In pg_upgrade, force timeline 1 in the new cluster (Bruce Momjian)

This change prevents upgrade failures caused by bogus complaints about missing WAL history files.

- In pg_upgrade, check for improperly non-connectable databases before proceeding (Bruce Momjian)
- In pg_upgrade, quote directory paths properly in the generated delete_old_cluster script (Bruce Momjian)
- In pg_upgrade, preserve database-level freezing info properly (Bruce Momjian)

This oversight could cause missing-clog-file errors for tables within the postgres and template1 databases.

- Run pg_upgrade and pg_resetxlog with restricted privileges on Windows, so that they don't fail when run by an administrator (Muhammad Asif Naeem)
- Improve handling of readdir() failures when scanning directories in initdb and pg_basebackup (Marco Nenciarini)
- Fix slow sorting algorithm in contrib/intarray (Tom Lane)
• Fix compile failure on Sparc V8 machines (Rob Rowan)

• Silence some build warnings on macOS (Tom Lane)

• Update time zone data files to tzdata release 2015d for DST law changes in Egypt, Mongolia, and Palestine, plus historical changes in Canada and Chile. Also adopt revised zone abbreviations for the America/Adak zone (HST/HDT not HAST/HADT).

E.78. Release 9.3.6

Release date: 2015-02-05

This release contains a variety of fixes from 9.3.5. For information about new features in the 9.3 major release, see Section E.84.

E.78.1. Migration to Version 9.3.6

A dump/restore is not required for those running 9.3.X. However, if you are a Windows user and are using the “Norwegian (Bokmål)” locale, manual action is needed after the upgrade to replace any “Norwegian (Bokmål)_Norway” locale names stored in PostgreSQL system catalogs with the plain-ASCII alias “Norwegian_Norway”. For details see http://wiki.postgresql.org/wiki/Changes_To_Norwegian_Locale

Also, if you are upgrading from a version earlier than 9.3.5, see Section E.79.

E.78.2. Changes

• Fix buffer overruns in to_char() (Bruce Momjian)

  When to_char() processes a numeric formatting template calling for a large number of digits, PostgreSQL would read past the end of a buffer. When processing a crafted timestamp formatting template, PostgreSQL would write past the end of a buffer. Either case could crash the server. We have not ruled out the possibility of attacks that lead to privilege escalation, though they seem unlikely. (CVE-2015-0241)

• Fix buffer overrun in replacement *printf() functions (Tom Lane)

  PostgreSQL includes a replacement implementation of printf and related functions. This code will overrun a stack buffer when formatting a floating point number (conversion specifiers e, E, f, F, g or G) with requested precision greater than about 500. This will crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. A database user can trigger such a buffer overrun through the to_char() SQL function. While that is the only affected core PostgreSQL functionality, extension modules that use printf-family functions may be at risk as well. This issue primarily affects PostgreSQL on Windows. PostgreSQL uses the system implementation of these functions where adequate, which it is on other modern platforms. (CVE-2015-0242)

• Fix buffer overruns in contrib/pgcrypto (Marko Tiikkaja, Noah Misch)

  Errors in memory size tracking within the pgcrypto module permitted stack buffer overruns and improper dependence on the contents of uninitialized memory. The buffer overrun cases can crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. (CVE-2015-0243)

• Fix possible loss of frontend/backend protocol synchronization after an error (Heikki Linnakangas)

  If any error occurred while the server was in the middle of reading a protocol message from the client, it could lose synchronization and incorrectly try to interpret part of the message's data as a new
protocol message. An attacker able to submit crafted binary data within a command parameter might succeed in injecting his own SQL commands this way. Statement timeout and query cancellation are the most likely sources of errors triggering this scenario. Particularly vulnerable are applications that use a timeout and also submit arbitrary user-crafted data as binary query parameters. Disabling statement timeout will reduce, but not eliminate, the risk of exploit. Our thanks to Emil Lenngren for reporting this issue. (CVE-2015-0244)

• Fix information leak via constraint-violation error messages (Stephen Frost)

Some server error messages show the values of columns that violate a constraint, such as a unique constraint. If the user does not have SELECT privilege on all columns of the table, this could mean exposing values that the user should not be able to see. Adjust the code so that values are displayed only when they came from the SQL command or could be selected by the user. (CVE-2014-8161)

• Lock down regression testing's temporary installations on Windows (Noah Misch)

Use SSPI authentication to allow connections only from the OS user who launched the test suite. This closes on Windows the same vulnerability previously closed on other platforms, namely that other users might be able to connect to the test postmaster. (CVE-2014-0067)

• Cope with the Windows locale named “Norwegian (Bokmål)” (Heikki Linnakangas)

Non-ASCII locale names are problematic since it's not clear what encoding they should be represented in. Map the troublesome locale name to a plain-ASCII alias, “Norwegian_Norway”.

• Avoid possible data corruption if ALTER DATABASE SET TABLESPACE is used to move a database to a new tablespace and then shortly later move it back to its original tablespace (Tom Lane)

• Avoid corrupting tables when ANALYZE inside a transaction is rolled back (Andres Freund, Tom Lane, Michael Paquier)

If the failing transaction had earlier removed the last index, rule, or trigger from the table, the table would be left in a corrupted state with the relevant pg_class flags not set though they should be.

• Ensure that unlogged tables are copied correctly during CREATE DATABASE or ALTER DATABASE SET TABLESPACE (Pavan Deolasee, Andres Freund)

• Fix incorrect processing of CreateEventTrigStmt.eventname (Petr Jelinek)

This could result in misbehavior if CREATE EVENT TRIGGER were executed as a prepared query, or via extended query protocol.

• Fix DROP’s dependency searching to correctly handle the case where a table column is recursively visited before its table (Petr Jelinek, Tom Lane)

This case is only known to arise when an extension creates both a datatype and a table using that datatype. The faulty code might refuse a DROP EXTENSION unless CASCADE is specified, which should not be required.

• Fix use-of-already-freed-memory problem in EvalPlanQual processing (Tom Lane)

In READ COMMITTED mode, queries that lock or update recently-updated rows could crash as a result of this bug.

• Avoid possible deadlock while trying to acquire tuple locks in EvalPlanQual processing (Álvaro Herrera, Mark Kirkwood)

• Fix failure to wait when a transaction tries to acquire a FOR NO KEY EXCLUSIVE tuple lock, while multiple other transactions currently hold FOR SHARE locks (Álvaro Herrera)
• Fix planning of SELECT FOR UPDATE when using a partial index on a child table (Kyotaro Horiguchi)

In READ COMMITTED mode, SELECT FOR UPDATE must also recheck the partial index's WHERE condition when rechecking a recently-updated row to see if it still satisfies the query's WHERE condition. This requirement was missed if the index belonged to an inheritance child table, so that it was possible to incorrectly return rows that no longer satisfy the query condition.

• Fix corner case wherein SELECT FOR UPDATE could return a row twice, and possibly miss returning other rows (Tom Lane)

In READ COMMITTED mode, a SELECT FOR UPDATE that is scanning an inheritance tree could incorrectly return a row from a prior child table instead of the one it should return from a later child table.

• Improve performance of EXPLAIN with large range tables (Tom Lane)

• Reject duplicate column names in the referenced-columns list of a FOREIGN KEY declaration (David Rowley)

This restriction is per SQL standard. Previously we did not reject the case explicitly, but later on the code would fail with bizarre-looking errors.

• Re-enable error for SELECT ... OFFSET -1 (Tom Lane)

A negative offset value has been an error since 8.4, but an optimization added in 9.3 accidentally turned the case into a no-op. Restore the expected behavior.

• Restore previous behavior of conversion of domains to JSON (Tom Lane)

This change causes domains over numeric and boolean to be treated like their base types for purposes of conversion to JSON. It worked like that before 9.3.5 and 9.2.9, but was unintentionally changed while fixing a related problem.

• Fix json_agg() to not return extra trailing right brackets in its result (Tom Lane)

• Fix bugs in raising a numeric value to a large integral power (Tom Lane)

The previous code could get a wrong answer, or consume excessive amounts of time and memory before realizing that the answer must overflow.

• In numeric_recv(), truncate away any fractional digits that would be hidden according to the value's dscale field (Tom Lane)

A numeric value's display scale (dscale) should never be less than the number of nonzero fractional digits; but apparently there's at least one broken client application that transmits binary numeric values in which that's true. This leads to strange behavior since the extra digits are taken into account by arithmetic operations even though they aren't printed. The least risky fix seems to be to truncate away such “hidden” digits on receipt, so that the value is indeed what it prints as.

• Fix incorrect search for shortest-first regular expression matches (Tom Lane)

Matching would often fail when the number of allowed iterations is limited by a ? quantifier or a bound expression.

• Reject out-of-range numeric timezone specifications (Tom Lane)

Simple numeric timezone specifications exceeding +/- 168 hours (one week) would be accepted, but could then cause null-pointer dereference crashes in certain operations. There's no use-case for such large UTC offsets, so reject them.

• Fix bugs in tsquery @> tsquery operator (Heikki Linnakangas)
Two different terms would be considered to match if they had the same CRC. Also, if the second operand had more terms than the first, it would be assumed not to be contained in the first; which is wrong since it might contain duplicate terms.

- Improve ispell dictionary’s defenses against bad affix files (Tom Lane)
- Allow more than 64K phrases in a thesaurus dictionary (David Boutin)

The previous coding could crash on an oversize dictionary, so this was deemed a back-patchable bug fix rather than a feature addition.

- Fix namespace handling in `xpath()` (Ali Akbar)

Previously, the `xml` value resulting from an `xpath()` call would not have namespace declarations if the namespace declarations were attached to an ancestor element in the input `xml` value, rather than to the specific element being returned. Propagate the ancestral declaration so that the result is correct when considered in isolation.

- Ensure that whole-row variables expose nonempty column names to functions that pay attention to column names within composite arguments (Tom Lane)

In some contexts, constructs like `row_to_json(tab.*)` may not produce the expected column names. This is fixed properly as of 9.4; in older branches, just ensure that we produce some nonempty name. (In some cases this will be the underlying table's column name rather than the query-assigned alias that should theoretically be visible.)

- Fix mishandling of system columns, particularly `tableoid`, in FDW queries (Etsuro Fujita)

- Fix assorted oversights in range-operator selectivity estimation (Emre Hasegeli)

This patch fixes corner-case “unexpected operator NNNN” planner errors, and improves the selectivity estimates for some other cases.

- Avoid doing `indexed_column = ANY (array)` as an index qualifier if that leads to an inferior plan (Andrew Gierth)

In some cases, `= ANY` conditions applied to non-first index columns would be done as index conditions even though it would be better to use them as simple filter conditions.

- Fix “variable not found in subplan target list” planner failure when an inline-able SQL function taking a composite argument is used in a `LATERAL` subselect and the composite argument is a lateral reference (Tom Lane)

- Fix planner problems with nested append relations, such as inherited tables within `UNION ALL` subqueries (Tom Lane)

- Fail cleanly when a GiST index tuple doesn't fit on a page, rather than going into infinite recursion (Andrew Gierth)

- Exempt tables that have per-table `cost_limit` and/or `cost_delay` settings from autovacuum's global cost balancing rules (Álvaro Herrera)

The previous behavior resulted in basically ignoring these per-table settings, which was unintended. Now, a table having such settings will be vacuumed using those settings, independently of what is going on in other autovacuum workers. This may result in heavier total I/O load than before, so such settings should be re-examined for sanity.

- Avoid wholesale autovacuuming when autovacuum is nominally off (Tom Lane)

Even when autovacuum is nominally off, we will still launch autovacuum worker processes to vacuum tables that are at risk of XID wraparound. However, such a worker process then proceeded.
to vacuum all tables in the target database, if they met the usual thresholds for autovacuuming. This is at best pretty unexpected; at worst it delays response to the wraparound threat. Fix it so that if autovacuum is turned off, workers only do anti-wraparound vacuums and not any other work.

• During crash recovery, ensure that unlogged relations are rewritten as empty and are synced to disk before recovery is considered complete (Abhijit Menon-Sen, Andres Freund)

This prevents scenarios in which unlogged relations might contain garbage data following database crash recovery.

• Fix race condition between hot standby queries and replaying a full-page image (Heikki Linnakangas)

This mistake could result in transient errors in queries being executed in hot standby.

• Fix several cases where recovery logic improperly ignored WAL records for COMMIT/ABORT PREPARED (Heikki Linnakangas)

The most notable oversight was that recovery_target_xid could not be used to stop at a two-phase commit.

• Prevent latest WAL file from being archived a second time at completion of crash recovery (Fujii Masao)

• Avoid creating unnecessary .ready marker files for timeline history files (Fujii Masao)

• Fix possible null pointer dereference when an empty prepared statement is used and the log_statement setting is mod or ddl (Fujii Masao)

• Change “pgstat wait timeout” warning message to be LOG level, and rephrase it to be more understandable (Tom Lane)

This message was originally thought to be essentially a can’t-happen case, but it occurs often enough on our slower buildfarm members to be a nuisance. Reduce it to LOG level, and expend a bit more effort on the wording: it now reads “using stale statistics instead of current ones because stats collector is not responding”.

• Fix possible corruption of postmaster’s list of dynamic background workers (Andres Freund)

• Fix SPARC spinlock implementation to ensure correctness if the CPU is being run in a non-TSO coherency mode, as some non-Solaris kernels do (Andres Freund)

• Warn if macOS's setlocale() starts an unwanted extra thread inside the postmaster (Noah Misch)

• Fix processing of repeated dbname parameters in PQconnectdbParams() (Alex Shulgin)

Unexpected behavior ensued if the first occurrence of dbname contained a connection string or URI to be expanded.

• Ensure that libpq reports a suitable error message on unexpected socket EOF (Marko Tiikkaja, Tom Lane)

Depending on kernel behavior, libpq might return an empty error string rather than something useful when the server unexpectedly closed the socket.

• Clear any old error message during PQreset() (Heikki Linnakangas)

If PQreset() is called repeatedly, and the connection cannot be re-established, error messages from the failed connection attempts kept accumulating in the PGconn’s error string.
• Properly handle out-of-memory conditions while parsing connection options in libpq (Alex Shulgin, Heikki Linnakangas)

• Fix array overrun in ecpg's version of ParseDateTime() (Michael Paquier)

• In initdb, give a clearer error message if a password file is specified but is empty (Mats Erik Andersson)

• Fix psql's \s command to work nicely with libedit, and add pager support (Stepan Rutz, Tom Lane)

  When using libedit rather than readline, \s printed the command history in a fairly unreadable encoded format, and on recent libedit versions might fail altogether. Fix that by printing the history ourselves rather than having the library do it. A pleasant side-effect is that the pager is used if appropriate.

  This patch also fixes a bug that caused newline encoding to be applied inconsistently when saving the command history with libedit. Multiline history entries written by older psql versions will be read cleanly with this patch, but perhaps not vice versa, depending on the exact libedit versions involved.

• Improve consistency of parsing of psql's special variables (Tom Lane)

  Allow variant spellings of on and off (such as 1/0) for ECHO_HIDDEN and ON_ERROR_ROLLBACK. Report a warning for unrecognized values for COMP_KEYWORD_CASE, ECHO, ECHO_HIDDEN, HISTCONTROL, ON_ERROR_ROLLBACK, and VERBOSITY. Recognize all values for all these variables case-insensitively; previously there was a mishmash of case-sensitive and case-insensitive behaviors.

• Make psql's \watch command display nulls as specified by \pset null (Fujii Masao)

• Fix psql's expanded-mode display to work consistently when using border = 3 and linestyle = ascii or unicode (Stephen Frost)

• Fix pg_dump to handle comments on event triggers without failing (Tom Lane)

• Allow parallel pg_dump to use --serializable-deferrable (Kevin Grittner)

• Improve performance of pg_dump when the database contains many instances of multiple dependency paths between the same two objects (Tom Lane)

• Fix pg_dumpall to restore its ability to dump from pre-8.1 servers (Gilles Darold)

• Fix possible deadlock during parallel restore of a schema-only dump (Robert Haas, Tom Lane)

• Fix core dump in pg_dump --binary-upgrade on zero-column composite type (Rushabh Lathia)

• Fix failure to fsync tables in nondefault tablespaces during pg_upgrade (Abhijit Menon-Sen, Andres Freund)

  With an operating system crash and some bad luck, this could result in data loss during an upgrade.

• In pg_upgrade, cope with cases where the new cluster creates a TOAST table for a table that didn't previously have one (Bruce Momjian)

  Previously this could result in failures due to OID conflicts.

• In pg_upgrade, don't try to set autovacuum_multixact_freeze_max_age for the old cluster (Bruce Momjian)

  This could result in failure because not all 9.3.X versions have that parameter. Fortunately, we don't actually need to set it at all.
Release Notes

- In pg_upgrade, preserve the transaction ID epoch (Bruce Momjian)
  This oversight did not bother PostgreSQL proper, but could confuse some external replication tools.
- Prevent WAL files created by `pg_basebackup -x/-X` from being archived again when the standby is promoted (Andres Freund)
- Fix memory leak in `pg_receivexlog` (Fujii Masao)
- Fix unintended suppression of `pg_receivexlog` verbose messages (Fujii Masao)
- Fix failure of `contrib/auto_explain` to print per-node timing information when doing `EXPLAIN ANALYZE` (Tom Lane)
- Fix upgrade-from-unpackaged script for `contrib/citext` (Tom Lane)
- Avoid integer overflow and buffer overrun in `contrib/hstore's hstore_to_json()` (Heikki Linnakangas)
- Fix recognition of numbers in `hstore_to_json_loose()`, so that JSON numbers and strings are correctly distinguished (Andrew Dunstan)
- Fix block number checking in `contrib/pageinspect's get_raw_page()` (Tom Lane)
  The incorrect checking logic could prevent access to some pages in non-main relation forks.
- Fix `contrib/pgcrypto's pgp_sym_decrypt()` to not fail on messages whose length is 6 less than a power of 2 (Marko Tiikkaja)
- Fix file descriptor leak in `contrib/pg_test_fsync` (Jeff Janes)
  This could cause failure to remove temporary files on Windows.
- Handle unexpected query results, especially NULLs, safely in `contrib/tablefunc's connectby()` (Michael Paquier)
  `connectby()` previously crashed if it encountered a NULL key value. It now prints that row but doesn't recurse further.
- Avoid a possible crash in `contrib/xml2's xslt_process()` (Mark Simonetti)
  `libxslt` seems to have an undocumented dependency on the order in which resources are freed; reorder our calls to avoid a crash.
- Mark some `contrib` I/O functions with correct volatility properties (Tom Lane)
  The previous over-conservative marking was immaterial in normal use, but could cause optimization problems or rejection of valid index expression definitions. Since the consequences are not large, we've just adjusted the function definitions in the extension modules' scripts, without changing version numbers.
- Numerous cleanups of warnings from Coverity static code analyzer (Andres Freund, Tatsuo Ishii, Marko Kreen, Tom Lane, Michael Paquier)
  These changes are mostly cosmetic but in some cases fix corner-case bugs, for example a crash rather than a proper error report after an out-of-memory failure. None are believed to represent security issues.
- Fix setup of background workers in EXEC_BACKEND builds, eg Windows (Robert Haas)
- Detect incompatible OpenLDAP versions during build (Noah Misch)
With OpenLDAP versions 2.4.24 through 2.4.31, inclusive, PostgreSQL backends can crash at exit. Raise a warning during configure based on the compile-time OpenLDAP version number, and test the crashing scenario in the contrib/dblink regression test.

- In non-MSVC Windows builds, ensure libpq.dll is installed with execute permissions (Noah Misch)

- Make pg_regress remove any temporary installation it created upon successful exit (Tom Lane)

  This results in a very substantial reduction in disk space usage during make check-world, since that sequence involves creation of numerous temporary installations.

- Support time zone abbreviations that change UTC offset from time to time (Tom Lane)

  Previously, PostgreSQL assumed that the UTC offset associated with a time zone abbreviation (such as EST) never changes in the usage of any particular locale. However this assumption fails in the real world, so introduce the ability for a zone abbreviation to represent a UTC offset that sometimes changes. Update the zone abbreviation definition files to make use of this feature in timezone locales that have changed the UTC offset of their abbreviations since 1970 (according to the IANA timezone database). In such timezones, PostgreSQL will now associate the correct UTC offset with the abbreviation depending on the given date.

- Update time zone abbreviations lists (Tom Lane)

  Add CST (China Standard Time) to our lists. Remove references to ADT as “Arabia Daylight Time”, an abbreviation that’s been out of use since 2007; therefore, claiming there is a conflict with “Atlantic Daylight Time” doesn’t seem especially helpful. Fix entirely incorrect GMT offsets for CKT (Cook Islands), FJT, and FJST (Fiji); we didn’t even have them on the proper side of the date line.

- Update time zone data files to tzdata release 2015a.

  The IANA timezone database has adopted abbreviations of the form A\times ST/A\times DT for all Australian time zones, reflecting what they believe to be current majority practice Down Under. These names do not conflict with usage elsewhere (other than ACST for Acre Summer Time, which has been in disuse since 1994). Accordingly, adopt these names into our “Default” timezone abbreviation set. The “Australia” abbreviation set now contains only CST, EAST, EST, SAST, SAT, and WST, all of which are thought to be mostly historical usage. Note that SAST has also been changed to be South Africa Standard Time in the “Default” abbreviation set.

  Also, add zone abbreviations SRET (Asia/Srednekolymsk) and XJT (Asia/Urumqi), and use WSST/WSDT for western Samoa. Also, there were DST law changes in Chile, Mexico, the Turks & Caicos Islands (America/Grand_Turk), and Fiji. There is a new zone Pacific/Bougainville for portions of Papua New Guinea. Also, numerous corrections for historical (pre-1970) time zone data.

### E.79. Release 9.3.5

**Release date:** 2014-07-24

This release contains a variety of fixes from 9.3.4. For information about new features in the 9.3 major release, see Section E.84.

### E.79.1. Migration to Version 9.3.5

A dump/restore is not required for those running 9.3.X.

However, this release corrects a logic error in pg_upgrade, as well as an index corruption problem in some GiST indexes. See the first two changelog entries below to find out whether your installation has been affected and what steps you should take if so.
E.79.2. Changes

- In pg_upgrade, remove pg_multixact files left behind by initdb (Bruce Momjian)

If you used a pre-9.3.5 version of pg_upgrade to upgrade a database cluster to 9.3, it might have left behind a file $PGDATA/pg_multixact/offsets/0000 that should not be there and will eventually cause problems in VACUUM. However, in common cases this file is actually valid and must not be removed. To determine whether your installation has this problem, run this query as superuser, in any database of the cluster:

```sql
WITH list(file) AS (SELECT * FROM pg_ls_dir('pg_multixact/offsets'))
SELECT EXISTS (SELECT * FROM list WHERE file = '0000') AND
NOT EXISTS (SELECT * FROM list WHERE file = '0001') AND
NOT EXISTS (SELECT * FROM list WHERE file = 'FFFF') AND
EXISTS (SELECT * FROM list WHERE file != '0000')
AS file_0000_removal_required;
```

If this query returns t, manually remove the file $PGDATA/pg_multixact/offsets/0000. Do nothing if the query returns f.

- Correctly initialize padding bytes in contrib/btree_gist indexes on bit columns (Heikki Linnakangas)

This error could result in incorrect query results due to values that should compare equal not being seen as equal. Users with GiST indexes on bit or bit varying columns should REINDEX those indexes after installing this update.

- Protect against torn pages when deleting GIN list pages (Heikki Linnakangas)

This fix prevents possible index corruption if a system crash occurs while the page update is being written to disk.

- Don't clear the right-link of a GiST index page while replaying updates from WAL (Heikki Linnakangas)

This error could lead to transiently wrong answers from GiST index scans performed in Hot Standby.

- Fix corner-case infinite loop during insertion into an SP-GiST text index (Tom Lane)

- Fix incorrect answers from SP-GiST index searches with −|− (range adjacency) operator (Heikki Linnakangas)

- Fix wraparound handling for pg_multixact/members (Álvaro Herrera)

- Truncate pg_multixact during checkpoints, not during VACUUM (Álvaro Herrera)

This change ensures that pg_multixact segments can't be removed if they'd still be needed during WAL replay after a crash.

- Fix possible inconsistency of all-visible flags after WAL recovery (Heikki Linnakangas)

- Fix possibly-incorrect cache invalidation during nested calls to ReceiveSharedInvalidMessages (Andres Freund)

- Fix race condition when updating a tuple concurrently locked by another process (Andres Freund, Álvaro Herrera)
• Fix “could not find pathkey item to sort” planner failures with \texttt{UNION ALL} over subqueries reading from tables with inheritance children (Tom Lane)

• Don’t assume a subquery’s output is unique if there’s a set-returning function in its targetlist (David Rowley)

  This oversight could lead to misoptimization of constructs like \texttt{WHERE \ x \ IN (SELECT y, generate_series(1,10) FROM t GROUP BY y)}.

• Improve planner to drop constant-NULL inputs of \texttt{AND/OR} when possible (Tom Lane)

  This change fixes some cases where the more aggressive parameter substitution done by 9.2 and later can lead to a worse plan than older versions produced.

• Ensure that the planner sees equivalent \texttt{VARIADIC} and non-\texttt{VARIADIC} function calls as equivalent (Tom Lane)

  This bug could for example result in failure to use expression indexes involving variadic functions. It might be necessary to re-create such indexes, and/or re-create views including variadic function calls that should match the indexes, for the fix to be effective for existing 9.3 installations.

• Fix handling of nested \texttt{JSON} objects in \texttt{json_populate_recordset()} and friends (Michael Paquier, Tom Lane)

  A nested \texttt{JSON} object could result in previous fields of the parent object not being shown in the output.

• Fix identification of input type category in \texttt{to_json()} and friends (Tom Lane)

  This is known to have led to inadequate quoting of \texttt{money} fields in the \texttt{JSON} result, and there may have been wrong results for other data types as well.

• Fix failure to detoast fields in composite elements of structured types (Tom Lane)

  This corrects cases where \texttt{TOAST} pointers could be copied into other tables without being dereferenced. If the original data is later deleted, it would lead to errors like “missing chunk number 0 for toast value ...” when the now-dangling pointer is used.

• Fix “record type has not been registered” failures with whole-row references to the output of \texttt{Append} plan nodes (Tom Lane)

• Fix possible crash when invoking a user-defined function while rewinding a cursor (Tom Lane)

• Fix query-lifespan memory leak while evaluating the arguments for a function in \texttt{FROM} (Tom Lane)

• Fix session-lifespan memory leaks in regular-expression processing (Tom Lane, Arthur O’Dwyer, Greg Stark)

• Fix data encoding error in \texttt{hungarian.stop} (Tom Lane)

• Prevent foreign tables from being created with OIDS when \texttt{default_with_oids} is true (Etsuro Fujita)

• Fix liveness checks for rows that were inserted in the current transaction and then deleted by a now-rolled-back subtransaction (Andres Freund)

  This could cause problems (at least spurious warnings, and at worst an infinite loop) if \texttt{CREATE INDEX} or \texttt{CLUSTER} were done later in the same transaction.

• Clear \texttt{pg_stat_activity.xact_start} during \texttt{PREPARE TRANSACTION} (Andres Freund)

  After the \texttt{PREPARE}, the originating session is no longer in a transaction, so it should not continue to display a transaction start time.
• Fix REASSIGN OWNED to not fail for text search objects (Álvaro Herrera)

• Prevent pg_class.relminmxid values from going backwards during VACUUM FULL (Álvaro Herrera)

• Reduce indentation in rule/view dumps to improve readability and avoid excessive whitespace (Greg Stark, Tom Lane)

This change reduces the amount of indentation applied to nested constructs, including some cases that the user probably doesn't think of as nested, such as UNION lists. Previously, deeply nested constructs were printed with an amount of whitespace growing as O(N^2), which created a performance problem and even risk of out-of-memory failures. Now the indentation is reduced modulo 40, which is initially odd to look at but seems to preserve readability better than simply limiting the indentation would do. Redundant parenthesization of UNION lists has been reduced as well.

• Fix dumping of rules/views when subsequent addition of a column has resulted in multiple input columns matching a USING specification (Tom Lane)

• Repair view printing for some cases involving functions in FROM that return a composite type containing dropped columns (Tom Lane)

• Block signals during postmaster startup (Tom Lane)

This ensures that the postmaster will properly clean up after itself if, for example, it receives SIGINT while still starting up.

• Fix client host name lookup when processing pg_hba.conf entries that specify host names instead of IP addresses (Tom Lane)

Ensure that reverse-DNS lookup failures are reported, instead of just silently not matching such entries. Also ensure that we make only one reverse-DNS lookup attempt per connection, not one per host name entry, which is what previously happened if the lookup attempts failed.

• Allow the root user to use postgres -C variable and postgres --describe-config (MauMau)

The prohibition on starting the server as root does not need to extend to these operations, and relaxing it prevents failure of pg_ctl in some scenarios.

• Secure Unix-domain sockets of temporary postmasters started during make check (Noah Misch)

Any local user able to access the socket file could connect as the server's bootstrap superuser, then proceed to execute arbitrary code as the operating-system user running the test, as we previously noted in CVE-2014-0067. This change defends against that risk by placing the server's socket in a temporary, mode 0700 subdirectory of /tmp. The hazard remains however on platforms where Unix sockets are not supported, notably Windows, because then the temporary postmaster must accept local TCP connections.

A useful side effect of this change is to simplify make check testing in builds that override DEFAULT_PGSOCKET_DIR. Popular non-default values like /var/run/postgresql are often not writable by the build user, requiring workarounds that will no longer be necessary.

• Fix tablespace creation WAL replay to work on Windows (MauMau)

• Fix detection of socket creation failures on Windows (Bruce Momjian)

• On Windows, allow new sessions to absorb values of PGC_BACKEND parameters (such as log_connections) from the configuration file (Amit Kapila)

Previously, if such a parameter were changed in the file post-startup, the change would have no effect.
Release Notes

- Properly quote executable path names on Windows (Nikhil Deshpande)
  
  This oversight could cause initdb and pg_upgrade to fail on Windows, if the installation path contained both spaces and @ signs.

- Fix linking of libpython on macOS (Tom Lane)
  
  The method we previously used can fail with the Python library supplied by Xcode 5.0 and later.

- Avoid buffer bloat in libpq when the server consistently sends data faster than the client can absorb it (Shin-ichi Morita, Tom Lane)
  
  libpq could be coerced into enlarging its input buffer until it runs out of memory (which would be reported misleadingly as “lost synchronization with server”). Under ordinary circumstances it’s quite far-fetched that data could be continuously transmitted more quickly than the recv() loop can absorb it, but this has been observed when the client is artificially slowed by scheduler constraints.

- Ensure that LDAP lookup attempts in libpq time out as intended (Laurenz Albe)

- Fix ecpg to do the right thing when an array of char * is the target for a FETCH statement returning more than one row, as well as some other array-handling fixes (Ashutosh Bapat)

- Fix pg_dump to cope with a materialized view that depends on a table's primary key (Tom Lane)
  
  This occurs if the view's query relies on functional dependency to abbreviate a GROUP BY list.
  
  pg_dump got sufficiently confused that it dumped the materialized view as a regular view.

- Fix parsing of pg_dumpall’s -i switch (Tom Lane)

- Fix pg_restore's processing of old-style large object comments (Tom Lane)
  
  A direct-to-database restore from an archive file generated by a pre-9.0 version of pg_dump would usually fail if the archive contained more than a few comments for large objects.

- Fix pg_upgrade for cases where the new server creates a TOAST table but the old version did not (Bruce Momjian)
  
  This rare situation would manifest as “relation OID mismatch” errors.

- In pg_upgrade, preserve pg_database.datminmxid and pg_class.relminmxid values from the old cluster, or insert reasonable values when upgrading from pre-9.3; also defend against unreasonable values in the core server (Bruce Momjian, Álvaro Herrera, Tom Lane)
  
  These changes prevent scenarios in which autovacuum might insist on scanning the entire cluster’s contents immediately upon starting the new cluster, or in which tracking of unfrozen MXID values might be disabled completely.

- Prevent contrib/auto_explain from changing the output of a user's EXPLAIN (Tom Lane)
  
  If auto_explain is active, it could cause an EXPLAIN (ANALYZE, TIMING OFF) command to nonetheless print timing information.

- Fix query-lifespan memory leak in contrib/dblink (MauMau, Joe Conway)

- In contrib/pgcrypto functions, ensure sensitive information is cleared from stack variables before returning (Marko Kreen)

- Prevent use of already-freed memory in contrib/pgstattuple's pgstat_heap() (Noah Misch)

- In contrib/uuid-oss, cache the state of the OSSP UUID library across calls (Tom Lane)
This improves the efficiency of UUID generation and reduces the amount of entropy drawn from /dev/urandom, on platforms that have that.

• Update time zone data files to tzdata release 2014e for DST law changes in Crimea, Egypt, and Morocco.

E.80. Release 9.3.4

Release date: 2014-03-20

This release contains a variety of fixes from 9.3.3. For information about new features in the 9.3 major release, see Section E.84.

E.80.1. Migration to Version 9.3.4

A dump/restore is not required for those running 9.3.X.

However, the error fixed in the first changelog entry below could have resulted in corrupt data on standby servers. It may be prudent to reinitialize standby servers from fresh base backups after installing this update.

Also, if you are upgrading from a version earlier than 9.3.3, see Section E.81.

E.80.2. Changes

• Fix WAL replay of locking an already-updated tuple (Andres Freund, Álvaro Herrera)

This error caused updated rows to not be found by index scans, resulting in inconsistent query results depending on whether an index scan was used. Subsequent processing could result in constraint violations, since the previously updated row would not be found by later index searches, thus possibly allowing conflicting rows to be inserted. Since this error is in WAL replay, it would only manifest during crash recovery or on standby servers. The improperly-replayed case most commonly arises when a table row that is referenced by a foreign-key constraint is updated concurrently with creation of a referencing row.

• Restore GIN metapages unconditionally to avoid torn-page risk (Heikki Linnakangas)

Although this oversight could theoretically result in a corrupted index, it is unlikely to have caused any problems in practice, since the active part of a GIN metapage is smaller than a standard 512-byte disk sector.

• Avoid race condition in checking transaction commit status during receipt of a NOTIFY message (Marko Tiikkaja)

This prevents a scenario wherein a sufficiently fast client might respond to a notification before database updates made by the notifier have become visible to the recipient.

• Allow materialized views to be referenced in UPDATE and DELETE commands (Michael Paquier)

Previously such queries failed with a complaint about not being able to lock rows in the materialized view.

• Allow regular-expression operators to be terminated early by query cancel requests (Tom Lane)

This prevents scenarios wherein a pathological regular expression could lock up a server process uninterruptibly for a long time.

• Remove incorrect code that tried to allow OVERLAPS with single-element row arguments (Joshua Yanovski)
This code never worked correctly, and since the case is neither specified by the SQL standard nor documented, it seemed better to remove it than fix it.

- Avoid getting more than AccessShareLock when de-parsing a rule or view (Dean Rasheed)

This oversight resulted in pg_dump unexpectedly acquiring RowExclusiveLock locks on tables mentioned as the targets of INSERT/UPDATE/DELETE commands in rules. While usually harmless, that could interfere with concurrent transactions that tried to acquire, for example, ShareLock on those tables.

- Improve performance of index endpoint probes during planning (Tom Lane)

This change fixes a significant performance problem that occurred when there were many not-yet-committed rows at the end of the index, which is a common situation for indexes on sequentially-assigned values such as timestamps or sequence-generated identifiers.

- Use non-default selectivity estimates for value IN (list) and value operator ANY (array) expressions when the righthand side is a stable expression (Tom Lane)

- Remove the correct per-database statistics file during DROP DATABASE (Tomas Vondra)

This fix prevents a permanent leak of statistics file space. Users who have done many DROP DATABASE commands since upgrading to PostgreSQL 9.3 may wish to check their statistics directory and delete statistics files that do not correspond to any existing database. Please note that db_0.stat should not be removed.

- Fix walsender ping logic to avoid inappropriate disconnects under continuous load (Andres Freund, Heikki Linnakangas)

walsender failed to send ping messages to the client if it was constantly busy sending WAL data; but it expected to see ping responses despite that, and would therefore disconnect once wal_sender_timeout elapsed.

- Fix walsender's failure to shut down cleanly when client is pg_recvverylog (Fujii Masao)

- Check WAL level and hot standby parameters correctly when doing crash recovery that will be followed by archive recovery (Heikki Linnakangas)

- Fix test to see if hot standby connections can be allowed immediately after a crash (Heikki Linnakangas)

- Add read-only data_checksums parameter to display whether page checksums are enabled (Heikki Linnakangas)

Without this parameter, determining the state of checksum processing was difficult.

- Prevent interrupts while reporting non-ERROR messages (Tom Lane)

This guards against rare server-process freezeups due to recursive entry to syslog(), and perhaps other related problems.

- Fix memory leak in PL/Perl when returning a composite result, including multiple-OUT-parameter cases (Alex Hunsaker)

- Fix tracking of psql script line numbers during \copy from out-of-line data (Kumar Rajeev Rastogi, Amit Khandekar)

\copy ... from incremented the script file line number for each data line, even if the data was not coming from the script file. This mistake resulted in wrong line numbers being reported for any errors occurring later in the same script file.
• Fix contrib/postgres_fdw to handle multiple join conditions properly (Tom Lane)

This oversight could result in sending WHERE clauses to the remote server for execution even though the clauses are not known to have the same semantics on the remote server (for example, clauses that use non-built-in operators). The query might succeed anyway, but it could also fail with errors from the remote server, or worse give silently wrong answers.

• Prevent intermittent “could not reserve shared memory region” failures on recent Windows versions (MauMau)

• Update time zone data files to tzdata release 2014a for DST law changes in Fiji and Turkey, plus historical changes in Israel and Ukraine.

E.81. Release 9.3.3

Release date: 2014-02-20

This release contains a variety of fixes from 9.3.2. For information about new features in the 9.3 major release, see Section E.84.

E.81.1. Migration to Version 9.3.3

A dump/restore is not required for those running 9.3.X.

However, several of the issues corrected in this release could have resulted in corruption of foreign-key constraints; that is, there might now be referencing rows for which there is no matching row in the referenced table. It may be worthwhile to recheck such constraints after installing this update. The simplest way to do that is to drop and recreate each suspect constraint; however, that will require taking an exclusive lock on both tables, so it is unlikely to be acceptable in production databases. Alternatively, you can do a manual join query between the two tables to look for unmatched rows.

Note also the requirement for replication standby servers to be upgraded before their master server is upgraded.

Also, if you are upgrading from a version earlier than 9.3.2, see Section E.82.

E.81.2. Changes

• Shore up GRANT ... WITH ADMIN OPTION restrictions (Noah Misch)

Granting a role without ADMIN OPTION is supposed to prevent the grantee from adding or removing members from the granted role, but this restriction was easily bypassed by doing SET ROLE first. The security impact is mostly that a role member can revoke the access of others, contrary to the wishes of his grantor. Unapproved role member additions are a lesser concern, since an uncooperative role member could provide most of his rights to others anyway by creating views or SECURITY DEFINER functions. (CVE-2014-0060)

• Prevent privilege escalation via manual calls to PL validator functions (Andres Freund)

The primary role of PL validator functions is to be called implicitly during CREATE FUNCTION, but they are also normal SQL functions that a user can call explicitly. Calling a validator on a function actually written in some other language was not checked for and could be exploited for privilege-escalation purposes. The fix involves adding a call to a privilege-checking function in each validator function. Non-core procedural languages will also need to make this change to their own validator functions, if any. (CVE-2014-0061)

• Avoid multiple name lookups during table and index DDL (Robert Haas, Andres Freund)

If the name lookups come to different conclusions due to concurrent activity, we might perform some parts of the DDL on a different table than other parts. At least in the case of CREATE INDEX,
this can be used to cause the permissions checks to be performed against a different table than the
index creation, allowing for a privilege escalation attack. (CVE-2014-0062)

• Prevent buffer overrun with long datetime strings (Noah Misch)

The maxdateLEN constant was too small for the longest possible value of type interval, al-
lowing a buffer overrun in interval_out(). Although the datetime input functions were more
careful about avoiding buffer overrun, the limit was short enough to cause them to reject some valid
inputs, such as input containing a very long timezone name. The ecpg library contained these vul-
nerabilities along with some of its own. (CVE-2014-0063)

• Prevent buffer overrun due to integer overflow in size calculations (Noah Misch, Heikki Linnakan-
gas)

Several functions, mostly type input functions, calculated an allocation size without checking for
overflow. If overflow did occur, a too-small buffer would be allocated and then written past.
(CVE-2014-0064)

• Prevent overruns of fixed-size buffers (Peter Eisentraut, Jozef Mlích)

Use strlcpy() and related functions to provide a clear guarantee that fixed-size buffers are not
overrun. Unlike the preceding items, it is unclear whether these cases really represent live issues,
since in most cases there appear to be previous constraints on the size of the input string. Nonetheless
it seems prudent to silence all Coverity warnings of this type. (CVE-2014-0065)

• Avoid crashing if crypt() returns NULL (Honza Horak, Bruce Momjian)

There are relatively few scenarios in which crypt() could return NULL, but contrib/chk-
pass would crash if it did. One practical case in which this could be an issue is if libc is configured
to refuse to execute unapproved hashing algorithms (e.g., “FIPS mode”). (CVE-2014-0066)

• Document risks of make check in the regression testing instructions (Noah Misch, Tom Lane)

Since the temporary server started by make check uses “trust” authentication, another user on the
same machine could connect to it as database superuser, and then potentially exploit the privileges of
the operating-system user who started the tests. A future release will probably incorporate changes
in the testing procedure to prevent this risk, but some public discussion is needed first. So for the
moment, just warn people against using make check when there are untrusted users on the same
machine. (CVE-2014-0067)

• Rework tuple freezing protocol (Álvaro Herrera, Andres Freund)

The logic for tuple freezing was unable to handle some cases involving freezing of multixact IDs,
with the practical effect that shared row-level locks might be forgotten once old enough.

Fixing this required changing the WAL record format for tuple freezing. While this is no issue
for standalone servers, when using replication it means that standby servers must be upgraded to
9.3.3 or later before their masters are. An older standby will be unable to interpret freeze records
generated by a newer master, and will fail with a PANIC message. (In such a case, upgrading the
standby should be sufficient to let it resume execution.)

• Create separate GUC parameters to control multixact freezing (Álvaro Herrera)

9.3 requires multixact tuple labels to be frozen before they grow too old, in the same fashion as
plain transaction ID labels have been frozen for some time. Previously, the transaction ID freezing
parameters were used for multixact IDs too; but since the consumption rates of transaction IDs
and multixact IDs can be quite different, this did not work very well. Introduce new settings vacuum_multixact_freeze_min_age, vacuum_multixact_freeze_table_age, and autovacuum_multixac-
t_freeze_max_age to control when to freeze multixacts.

• Account for remote row locks propagated by local updates (Álvaro Herrera)
If a row was locked by transaction A, and transaction B updated it, the new version of the row created by B would be locked by A, yet visible only to B. If transaction B then again updated the row, A's lock wouldn't get checked, thus possibly allowing B to complete when it shouldn't. This case is new in 9.3 since prior versions did not have any types of row locking that would permit another transaction to update the row at all.

This oversight could allow referential integrity checks to give false positives (for instance, allow deletes that should have been rejected). Applications using the new commands `SELECT FOR KEY SHARE` and `SELECT FOR NO KEY UPDATE` might also have suffered locking failures of this kind.

- Prevent “forgetting” valid row locks when one of several holders of a row lock aborts (Álvaro Herrera)

This was yet another mechanism by which a shared row lock could be lost, thus possibly allowing updates that should have been prevented by foreign-key constraints.

- Fix incorrect logic during update chain locking (Álvaro Herrera)

This mistake could result in spurious “could not serialize access due to concurrent update” errors in `REPEATABLE READ` and `SERIALIZABLE` transaction isolation modes.

- Handle wraparound correctly during extension or truncation of `pg_multixact/members` (Andres Freund, Álvaro Herrera)

- Fix handling of 5-digit filenames in `pg_multixact/members` (Álvaro Herrera)

As of 9.3, these names can be more than 4 digits, but the directory cleanup code ignored such files.

- Improve performance of multixact cache code (Álvaro Herrera)

- Optimize updating a row that’s already locked by the same transaction (Andres Freund, Álvaro Herrera)

This fixes a performance regression from pre-9.3 versions when doing `SELECT FOR UPDATE` followed by `UPDATE/DELETE`.

- During archive recovery, prefer highest timeline number when WAL segments with the same ID are present in both the archive and `pg_xlog/` (Kyotaro Horiguchi)

Previously, not-yet-archived segments could get ignored during recovery. This reverts an undesirable behavioral change in 9.3.0 back to the way things worked pre-9.3.

- Fix possible mis-replay of WAL records when some segments of a relation aren't full size (Greg Stark, Tom Lane)

The WAL update could be applied to the wrong page, potentially many pages past where it should have been. Aside from corrupting data, this error has been observed to result in significant “bloat” of standby servers compared to their masters, due to updates being applied far beyond where the end-of-file should have been. This failure mode does not appear to be a significant risk during crash recovery, only when initially synchronizing a standby created from a base backup taken from a quickly-changing master.

- Fix bug in determining when recovery has reached consistency (Tomonari Katsumata, Heikki Linnakangas)

In some cases WAL replay would mistakenly conclude that the database was already consistent at the start of replay, thus possibly allowing hot-standby queries before the database was really consistent. Other symptoms such as “PANIC: WAL contains references to invalid pages” were also possible.
• Fix WAL logging of visibility map changes (Heikki Linnakangas)

• Fix improper locking of btree index pages while replaying a VACUUM operation in hot-standby mode (Andres Freund, Heikki Linnakangas, Tom Lane)

  This error could result in “PANIC: WAL contains references to invalid pages” failures.

• Ensure that insertions into non-leaf GIN index pages write a full-page WAL record when appropriate (Heikki Linnakangas)

  The previous coding risked index corruption in the event of a partial-page write during a system crash.

• When `pause_at_recovery_target` and `recovery_target_inclusive` are both set, ensure the target record is applied before pausing, not after (Heikki Linnakangas)

• Ensure `walreceiver` sends hot-standby feedback messages on time even when there is a continuous stream of data (Andres Freund, Amit Kapila)

• Prevent timeout interrupts from taking control away from mainline code unless `ImmediateInterruptOK` is set (Andres Freund, Tom Lane)

  This is a serious issue for any application making use of statement timeouts, as it could cause all manner of strange failures after a timeout occurred. We have seen reports of “stuck” spinlocks, ERRORs being unexpectedly promoted to PANICs, unkillable backends, and other misbehaviors.

• Fix race conditions during server process exit (Robert Haas)

  Ensure that signal handlers don’t attempt to use the process's `MyProc` pointer after it’s no longer valid.

• Fix race conditions in `walsender` shutdown logic and `walreceiver` SIGHUP signal handler (Tom Lane)

• Fix unsafe references to `errno` within error reporting logic (Christian Kruse)

  This would typically lead to odd behaviors such as missing or inappropriate `HINT` fields.

• Fix possible crashes from using `ereport()` too early during server startup (Tom Lane)

  The principal case we've seen in the field is a crash if the server is started in a directory it doesn't have permission to read.

• Clear retry flags properly in OpenSSL socket write function (Alexander Kukushkin)

  This omission could result in a server lockup after unexpected loss of an SSL-encrypted connection.

• Fix length checking for Unicode identifiers (U&"..." syntax) containing escapes (Tom Lane)

  A spurious truncation warning would be printed for such identifiers if the escaped form of the identifier was too long, but the identifier actually didn't need truncation after de-escaping.

• Fix parsing of Unicode literals and identifiers just before the end of a command string or function body (Tom Lane)

• Allow keywords that are type names to be used in lists of roles (Stephen Frost)

  A previous patch allowed such keywords to be used without quoting in places such as role identifiers; but it missed cases where a list of role identifiers was permitted, such as DROP ROLE.

• Fix parser crash for `EXISTS(SELECT * FROM zero_column_table)` (Tom Lane)
• Fix possible crash due to invalid plan for nested sub-selects, such as WHERE (... × IN (SELECT ...) ... IN (SELECT ...) (Tom Lane)

• Fix mishandling of WHERE conditions pulled up from a LATERAL subquery (Tom Lane)

  The typical symptom of this bug was a “JOIN qualification cannot refer to other relations” error, though subtle logic errors in created plans seem possible as well.

• Disallow LATERAL references to the target table of an UPDATE/DELETE (Tom Lane)

  While this might be allowed in some future release, it was unintentional in 9.3, and didn’t work quite right anyway.

• Fix UPDATE/DELETE of an inherited target table that has UNION ALL subqueries (Tom Lane)

  Without this fix, UNION ALL subqueries aren’t correctly inserted into the update plans for inheritance child tables after the first one, typically resulting in no update happening for those child table(s).

• Fix ANALYZE to not fail on a column that’s a domain over a range type (Tom Lane)

• Ensure that ANALYZE creates statistics for a table column even when all the values in it are “too wide” (Tom Lane)

  ANALYZE intentionally omits very wide values from its histogram and most-common-values calculations, but it neglected to do something sane in the case that all the sampled entries are too wide.

• In ALTER TABLE ... SET TABLESPACE, allow the database’s default tablespace to be used without a permissions check (Stephen Frost)

  CREATE TABLE has always allowed such usage, but ALTER TABLE didn’t get the memo.

• Fix support for extensions containing event triggers (Tom Lane)

• Fix “cannot accept a set” error when some arms of a CASE return a set and others don’t (Tom Lane)

• Fix memory leakage in JSON functions (Craig Ringer)

• Properly distinguish numbers from non-numbers when generating JSON output (Andrew Dunstan)

• Fix checks for all-zero client addresses in pgstat functions (Kevin Grittner)

• Fix possible misclassification of multibyte characters by the text search parser (Tom Lane)

  Non-ASCII characters could be misclassified when using C locale with a multibyte encoding. On Cygwin, non-C locales could fail as well.

• Fix possible misbehavior in plainto_tquery() (Heikki Linnakangas)

  Use memmove() not memcpy() for copying overlapping memory regions. There have been no field reports of this actually causing trouble, but it’s certainly risky.

• Fix placement of permissions checks in pg_start_backup() and pg_stop_backup() (Andres Freund, Magnus Hagander)

  The previous coding might attempt to do catalog access when it shouldn’t.

• Accept SHIFT_JIS as an encoding name for locale checking purposes (Tatsuo Ishii)

• Fix *-qualification of named parameters in SQL-language functions (Tom Lane)

  Given a composite-type parameter named foo, $1.* worked fine, but foo.* not so much.
• Fix misbehavior of PQhost() on Windows (Fujii Masao)
  It should return localhost if no host has been specified.

• Improve error handling in libpq and psql for failures during COPY TO STDOUT/FROM STDIN
  (Tom Lane)
  In particular this fixes an infinite loop that could occur in 9.2 and up if the server connection was
  lost during COPY FROM STDIN. Variants of that scenario might be possible in older versions,
  or with other client applications.

• Fix incorrect translation handling in some psql \d commands (Peter Eisentraut, Tom Lane)

• Ensure pg_basebackup's background process is killed when exiting its foreground process (Magnus
  Hagander)

• Fix possible incorrect printing of filenames in pg_basebackup's verbose mode (Magnus Hagander)

• Avoid including tablespaces inside PGDATA twice in base backups (Dimitri Fontaine, Magnus
  Hagander)

• Fix misaligned descriptors in ecpg (MauMau)

• In ecpg, handle lack of a hostname in the connection parameters properly (Michael Meskes)

• Fix performance regression in contrib/dblink connection startup (Joe Conway)
  Avoid an unnecessary round trip when client and server encodings match.

• In contrib/isn, fix incorrect calculation of the check digit for ISMN values (Fabien Coelho)

• Fix contrib/pgbench's progress logging to avoid overflow when the scale factor is large (Tatsu-
  suo Ishii)

• Fix contrib/pg_stat_statement's handling of CURRENT_DATE and related constructs
  (Kyotaro Horiguchi)

• Improve lost-connection error handling in contrib/postgres_fdw (Tom Lane)

• Ensure client-code-only installation procedure works as documented (Peter Eisentraut)

• In Mingw and Cygwin builds, install the libpq DLL in the bin directory (Andrew Dunstan)
  This duplicates what the MSVC build has long done. It should fix problems with programs like psql
  failing to start because they can't find the DLL.

• Avoid using the deprecated dllwrap tool in Cygwin builds (Marco Atzeri)

• Enable building with Visual Studio 2013 (Brar Piening)

• Don't generate plain-text HISTORY and src/test/regress/README files anymore (Tom
  Lane)
  These text files duplicated the main HTML and PDF documentation formats. The trouble involved
  in maintaining them greatly outweighs the likely audience for plain-text format. Distribution tarballs
  will still contain files by these names, but they'll just be stubs directing the reader to consult the
  main documentation. The plain-text INSTALL file will still be maintained, as there is arguably a
  use-case for that.

• Update time zone data files to tzdata release 2013i for DST law changes in Jordan and historical
  changes in Cuba.
In addition, the zones Asia/Riyadh87, Asia/Riyadh88, and Asia/Riyadh89 have been removed, as they are no longer maintained by IANA, and never represented actual civil timekeeping practice.

E.82. Release 9.3.2

Release date: 2013-12-05

This release contains a variety of fixes from 9.3.1. For information about new features in the 9.3 major release, see Section E.84.

E.82.1. Migration to Version 9.3.2

A dump/restore is not required for those running 9.3.X.

However, this release corrects a number of potential data corruption issues. See the first three changelog entries below to find out whether your installation has been affected and what steps you can take if so.

Also, if you are upgrading from a version earlier than 9.3.1, see Section E.83.

E.82.2. Changes

• Fix VACUUM’s tests to see whether it can update relfrozenxid (Andres Freund)

    In some cases VACUUM (either manual or autovacuum) could incorrectly advance a table’s relfrozenxid value, allowing tuples to escape freezing, causing those rows to become invisible once 2^31 transactions have elapsed. The probability of data loss is fairly low since multiple incorrect advancements would need to happen before actual loss occurs, but it’s not zero. In 9.2.0 and later, the probability of loss is higher, and it’s also possible to get “could not access status of transaction” errors as a consequence of this bug. Users upgrading from releases 9.0.4 or 8.4.8 or earlier are not affected, but all later versions contain the bug.

    The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix latent corruption but will not be able to fix all pre-existing data errors. However, an installation can be presumed safe after performing this vacuuming if it has executed fewer than 2^31 update transactions in its lifetime (check this with SELECT txid_current() < 2^31).

• Fix multiple bugs in MultiXactId freezing (Andres Freund, Álvaro Herrera)

    These bugs could lead to “could not access status of transaction” errors, or to duplicate or vanishing rows. Users upgrading from releases prior to 9.3.0 are not affected.

    The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix latent corruption but will not be able to fix all pre-existing data errors.

    As a separate issue, these bugs can also cause standby servers to get out of sync with the primary, thus exhibiting data errors that are not in the primary. Therefore, it’s recommended that 9.3.0 and 9.3.1 standby servers be re-cloned from the primary (e.g., with a new base backup) after upgrading.

• Fix initialization of pg_clog and pg_subtrans during hot standby startup (Andres Freund, Heikki Linnakangas)

    This bug can cause data loss on standby servers at the moment they start to accept hot-standby queries, by marking committed transactions as uncommitted. The likelihood of such corruption is
small unless, at the time of standby startup, the primary server has executed many updating transactions since its last checkpoint. Symptoms include missing rows, rows that should have been deleted being still visible, and obsolete versions of updated rows being still visible alongside their newer versions.

This bug was introduced in versions 9.3.0, 9.2.5, 9.1.10, and 9.0.14. Standby servers that have only been running earlier releases are not at risk. It's recommended that standby servers that have ever run any of the buggy releases be re-cloned from the primary (e.g., with a new base backup) after upgrading.

- Fix multiple bugs in update chain traversal (Andres Freund, Álvaro Herrera)
  These bugs could result in incorrect behavior, such as locking or even updating the wrong row, in the presence of concurrent updates. Spurious “unable to fetch updated version of tuple” errors were also possible.

- Fix dangling-pointer problem in fast-path locking (Tom Lane)
  This could lead to corruption of the lock data structures in shared memory, causing “lock already held” and other odd errors.

- Fix assorted race conditions in timeout management (Tom Lane)
  These errors could result in a server process becoming unresponsive because it had blocked SIGALRM and/or SIGINT.

- Truncate pg_multixact contents during WAL replay (Andres Freund)
  This avoids ever-increasing disk space consumption in standby servers.

- Ensure an anti-wraparound VACUUM counts a page as scanned when it's only verified that no tuples need freezing (Sergey Burladyan, Jeff Janes)
  This bug could result in failing to advance relfrozenxid, so that the table would still be thought to need another anti-wraparound vacuum. In the worst case the database might even shut down to prevent wraparound.

- Fix full-table-vacuum request mechanism for MultiXactIds (Andres Freund)
  This bug could result in large amounts of useless autovacuum activity.

- Fix race condition in GIN index posting tree page deletion (Heikki Linnakangas)
  This could lead to transient wrong answers or query failures.

- Fix “unexpected spgdoinser() failure” error during SP-GiST index creation (Teodor Sigaev)

- Fix assorted bugs in materialized views (Kevin Grittner, Andres Freund)

- Re-allow duplicate table aliases if they're within aliased JOINs (Tom Lane)
  Historically PostgreSQL has accepted queries like
  
  ```sql
  SELECT ... FROM tab1 x CROSS JOIN (tab2 x CROSS JOIN tab3 y) z
  ```
  although a strict reading of the SQL standard would forbid the duplicate usage of table alias x. A misguided change in 9.3.0 caused it to reject some such cases that were formerly accepted. Restore the previous behavior.

- Avoid flattening a subquery whose SELECT list contains a volatile function wrapped inside a sub-SELECT (Tom Lane)
This avoids unexpected results due to extra evaluations of the volatile function.

- Fix planner's processing of non-simple-variable subquery outputs nested within outer joins (Tom Lane)

  This error could lead to incorrect plans for queries involving multiple levels of subqueries within JOIN syntax.

- Fix incorrect planning in cases where the same non-strict expression appears in multiple WHERE and outer JOIN equality clauses (Tom Lane)

- Fix planner crash with whole-row reference to a subquery (Tom Lane)

- Fix incorrect generation of optimized MIN()/MAX() plans for inheritance trees (Tom Lane)

  The planner could fail in cases where the MIN()/MAX() argument was an expression rather than a simple variable.

- Fix premature deletion of temporary files (Andres Freund)

- Prevent intra-transaction memory leak when printing range values (Tom Lane)

  This fix actually cures transient memory leaks in any datatype output function, but range types are the only ones known to have had a significant problem.

- Fix memory leaks when reloading configuration files (Heikki Linnakangas, Hari Babu)

- Prevent incorrect display of dropped columns in NOT NULL and CHECK constraint violation messages (Michael Paquier and Tom Lane)

- Allow default arguments and named-argument notation for window functions (Tom Lane)

  Previously, these cases were likely to crash.

- Suppress trailing whitespace on each line when pretty-printing rules and views (Tom Lane)

  9.3.0 generated such whitespace in many more cases than previous versions did. To reduce unexpected behavioral changes, suppress unnecessary whitespace in all cases.

- Fix possible read past end of memory in rule printing (Peter Eisentraut)

- Fix array slicing of int2vector and oidvector values (Tom Lane)

  Expressions of this kind are now implicitly promoted to regular int2 or oid arrays.

- Return a valid JSON value when converting an empty hstore value to json (Oskari Saarenmaa)

- Fix incorrect behaviors when using a SQL-standard, simple GMT offset timezone (Tom Lane)

  In some cases, the system would use the simple GMT offset value when it should have used the regular timezone setting that had prevailed before the simple offset was selected. This change also causes the timeofday function to honor the simple GMT offset zone.

- Prevent possible misbehavior when logging translations of Windows error codes (Tom Lane)

- Properly quote generated command lines in pg_ctl (Naoya Anzai and Tom Lane)

  This fix applies only to Windows.

- Fix pg_dumpall to work when a source database sets default_transaction_read_only via ALTER DATABASE SET (Kevin Grittner)
Previously, the generated script would fail during restore.

- Fix pg_isready to handle its `-d` option properly (Fabrízio de Royes Mello and Fujii Masao)
- Fix parsing of WAL file names in pg_receivexlog (Heikki Linnakangas)
  This error made pg_receivexlog unable to restart streaming after stopping, once at least 4 GB of WAL had been written.
- Report out-of-disk-space failures properly in pg_upgrade (Peter Eisentraut)
- Make ecpg search for quoted cursor names case-sensitively (Zoltán Böszörményi)
- Fix ecpg's processing of lists of variables declared varchar (Zoltán Böszörményi)
- Make contrib/lo defend against incorrect trigger definitions (Marc Cousin)
- Update time zone data files to tzdata release 2013h for DST law changes in Argentina, Brazil, Jordan, Libya, Liechtenstein, Morocco, and Palestine. Also, new timezone abbreviations WIB, WIT, WITA for Indonesia.

**E.83. Release 9.3.1**

**Release date:** 2013-10-10

This release contains a variety of fixes from 9.3.0. For information about new features in the 9.3 major release, see Section E.84.

**E.83.1. Migration to Version 9.3.1**

A dump/restore is not required for those running 9.3.X.

However, if you use the hstore extension, see the first changelog entry.

**E.83.2. Changes**

- Ensure new-in-9.3 JSON functionality is added to the hstore extension during an update (Andrew Dunstan)

  Users who upgraded a pre-9.3 database containing hstore should execute

  ```sql
  ALTER EXTENSION hstore UPDATE;
  ```

  after installing 9.3.1, to add two new JSON functions and a cast. (If hstore is already up to date, this command does nothing.)

- Fix memory leak when creating B-tree indexes on range columns (Heikki Linnakangas)
- Fix memory leak caused by lo_open() failure (Heikki Linnakangas)
- Serializable snapshot fixes (Kevin Grittner, Heikki Linnakangas)
- Fix deadlock bug in libpq when using SSL (Stephen Frost)
- Fix timeline handling bugs in pg_receivexlog (Heikki Linnakangas, Andrew Gierth)
- Prevent CREATE FUNCTION from checking SET variables unless function body checking is enabled (Tom Lane)
• Remove rare inaccurate warning during vacuum of index-less tables (Heikki Linnakangas)

E.84. Release 9.3

Release date: 2013-09-09

E.84.1. Overview

Major enhancements in PostgreSQL 9.3 include:

• Add materialized views
• Make simple views auto-updatable
• Add many features for the JSON data type, including operators and functions to extract elements from JSON values
• Implement SQL-standard LATERAL option for FROM-clause subqueries and function calls
• Allow foreign data wrappers to support writes (inserts/updates/deletes) on foreign tables
• Add a Postgres foreign data wrapper to allow access to other Postgres servers
• Add support for event triggers
• Add optional ability to checksum data pages and report corruption
• Prevent non-key-field row updates from blocking foreign key checks
• Greatly reduce System V shared memory requirements

The above items are explained in more detail in the sections below.

E.84.2. Migration to Version 9.3

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.3 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

E.84.2.1. Server Settings

• Rename replication_timeout to wal_sender_timeout (Amit Kapila)

This setting controls the WAL sender timeout.

• Require superuser privileges to set commit_delay because it can now potentially delay other sessions (Simon Riggs)

• Allow in-memory sorts to use their full memory allocation (Jeff Janes)

Users who have set work_mem based on the previous behavior may need to revisit that setting.

E.84.2.2. Other

• Throw an error if a tuple to be updated or deleted has already been updated or deleted by a BEFORE trigger (Kevin Grittner)

Formerly, the originally-intended update was silently skipped, resulting in logical inconsistency since the trigger might have propagated data to other places based on the intended update. Now an
error is thrown to prevent the inconsistent results from being committed. If this change affects your application, the best solution is usually to move the data-propagation actions to an **AFTER** trigger.

This error will also be thrown if a query invokes a volatile function that modifies rows that are later modified by the query itself. Such cases likewise previously resulted in silently skipping updates.

- **Change multicolumn ON UPDATE SET NULL/SET DEFAULT foreign key actions to affect all columns of the constraint, not just those changed in the UPDATE** (Tom Lane)

  Previously, we would set only those referencing columns that correspond to referenced columns that were changed by the **UPDATE**. This was what was required by SQL-92, but more recent editions of the SQL standard specify the new behavior.

- **Force cached plans to be replanned if the search_path changes** (Tom Lane)

  Previously, cached plans already generated in the current session were not redone if the query was re-executed with a new **search_path** setting, resulting in surprising behavior.

- **Fix to_number() to properly handle a period used as a thousands separator** (Tom Lane)

  Previously, a period was considered to be a decimal point even when the locale says it isn’t and the D format code is used to specify use of the locale-specific decimal point. This resulted in wrong answers if FM format was also used.

- **Fix STRICT non-set-returning functions that have set-returning functions in their arguments to properly return null rows** (Tom Lane)

  A null value passed to the strict function should result in a null output, but instead, that output row was suppressed entirely.

- **Store WAL in a continuous stream, rather than skipping the last 16MB segment every 4GB** (Heikki Linnakangas)

  Previously, WAL files with names ending in FF were not used because of this skipping. If you have WAL backup or restore scripts that took this behavior into account, they will need to be adjusted.

- **In pg_constraint.confmatchtype, store the default foreign key match type (non-FULL, non-PARTIAL) as s for “simple”** (Tom Lane)

  Previously this case was represented by u for “unspecified”.

### E.84.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.3 and the previous major release.

#### E.84.3.1. Server

#### E.84.3.1.1. Locking

- **Prevent non-key-field row updates from blocking foreign key checks** (Álvaro Herrera, Noah Misch, Andres Freund, Alexander Shulgin, Marti Raudsepp, Alexander Shulgin)

  This change improves concurrency and reduces the probability of deadlocks when updating tables involved in a foreign-key constraint. **UPDATEs** that do not change any columns referenced in a foreign key now take the new **NO KEY UPDATE** lock mode on the row, while foreign key checks use the new **KEY SHARE** lock mode, which does not conflict with **NO KEY UPDATE**. So there is no blocking unless a foreign-key column is changed.

- **Add configuration variable lock_timeout to allow limiting how long a session will wait to acquire any one lock** (Zoltán Bőszörményi)
E.84.3.1.2. Indexes

- Add SP-GiST support for range data types (Alexander Korotkov)
- Allow GiST indexes to be unlogged (Jeevan Chalke)
- Improve performance of GiST index insertion by randomizing the choice of which page to descend to when there are multiple equally good alternatives (Heikki Linnakangas)
- Improve concurrency of hash index operations (Robert Haas)

E.84.3.1.3. Optimizer

- Collect and use histograms of upper and lower bounds, as well as range lengths, for range types (Alexander Korotkov)
- Improve optimizer's cost estimation for index access (Tom Lane)
- Improve optimizer's hash table size estimate for doing DISTINCT via hash aggregation (Tom Lane)
- Suppress no-op Result and Limit plan nodes (Kyotaro Horiguchi, Amit Kapila, Tom Lane)
- Reduce optimizer overhead by not keeping plans on the basis of cheap startup cost when the optimizer only cares about total cost overall (Tom Lane)

E.84.3.1.4. General Performance

- Add COPY FREEZE option to avoid the overhead of marking tuples as frozen later (Simon Riggs, Jeff Davis)
- Improve performance of NUMERIC calculations (Kyotaro Horiguchi)
- Improve synchronization of sessions waiting for commit_delay (Peter Geoghegan)
  This greatly improves the usefulness of commit_delay.
- Improve performance of the CREATE TEMPORARY TABLE ... ON COMMIT DELETE ROWS option by not truncating such temporary tables in transactions that haven't touched any temporary tables (Heikki Linnakangas)
- Make vacuum recheck visibility after it has removed expired tuples (Pavan Deolasee)
  This increases the chance of a page being marked as all-visible.
- Add per-resource-owner lock caches (Jeff Janes)
  This speeds up lock bookkeeping at statement completion in multi-statement transactions that hold many locks; it is particularly useful for pg_dump.
- Avoid scanning the entire relation cache at commit of a transaction that creates a new relation (Jeff Janes)
  This speeds up sessions that create many tables in successive small transactions, such as a pg_restore run.
- Improve performance of transactions that drop many relations (Tomas Vondra)

E.84.3.1.5. Monitoring

- Add optional ability to checksum data pages and report corruption (Simon Riggs, Jeff Davis, Greg Smith, Ants Aasma)
The checksum option can be set during initdb.

- Split the statistics collector's data file into separate global and per-database files (Tomas Vondra)
  This reduces the I/O required for statistics tracking.
- Fix the statistics collector to operate properly in cases where the system clock goes backwards (Tom Lane)
  Previously, statistics collection would stop until the time again reached the latest time previously recorded.
- Emit an informative message to postmaster standard error when we are about to stop logging there (Tom Lane)
  This should help reduce user confusion about where to look for log output in common configurations that log to standard error only during postmaster startup.

**E.84.3.1.6. Authentication**

- When an authentication failure occurs, log the relevant `pg_hba.conf` line, to ease debugging of unintended failures (Magnus Hagander)
- Improve LDAP error reporting and documentation (Peter Eisentraut)
- Add support for specifying LDAP authentication parameters in URL format, per RFC 4516 (Peter Eisentraut)
- Change the `ssl_ciphers` parameter to start with `DEFAULT`, rather than `ALL`, then remove insecure ciphers (Magnus Hagander)
  This should yield a more appropriate SSL cipher set.
- Parse and load `pg_ident.conf` once, not during each connection (Amit Kapila)
  This is similar to how `pg_hba.conf` is processed.

**E.84.3.1.7. Server Settings**

- Greatly reduce System V shared memory requirements (Robert Haas)
  On Unix-like systems, `mmap()` is now used for most of PostgreSQL’s shared memory. For most users, this will eliminate any need to adjust kernel parameters for shared memory.
- Allow the postmaster to listen on multiple Unix-domain sockets (Honza Horák)
  The configuration parameter `unix_socket_directory` is replaced by `unix_socket_directories`, which accepts a list of directories.
- Allow a directory of configuration files to be processed (Magnus Hagander, Greg Smith, Selena Deckelmann)
  Such a directory is specified with `include_dir` in the server configuration file.
- Increase the maximum initdb-configured value for `shared_buffers` to 128MB (Robert Haas)
  This is the maximum value that initdb will attempt to set in `postgresql.conf`; the previous maximum was 32MB.
- Remove the external PID file, if any, on postmaster exit (Peter Eisentraut)
E.84.3.2. Replication and Recovery

- Allow a streaming replication standby to follow a timeline switch (Heikki Linnakangas)

  This allows streaming standby servers to receive WAL data from a slave newly promoted to master status. Previously, other standbys would require a resync to begin following the new master.

- Add SQL functions `pg_is_in_backup()` and `pg_backup_start_time()` (Gilles Darold)

  These functions report the status of base backups.

- Improve performance of streaming log shipping with `synchronous_commit` disabled (Andres Freund)

- Allow much faster promotion of a streaming standby to primary (Simon Riggs, Kyotaro Horiguchi)

- Add the last checkpoint's redo location to `pg_controldata`'s output (Fujii Masao)

  This information is useful for determining which WAL files are needed for restore.

- Allow tools like `pg_receivexlog` to run on computers with different architectures (Heikki Linnakangas)

  WAL files can still only be replayed on servers with the same architecture as the primary; but they can now be transmitted to and stored on machines of any architecture, since the streaming replication protocol is now machine-independent.

- Make `pg_basebackup --write-recovery-conf` output a minimal `recovery.conf` file (Zoltán Böszörményi, Magnus Hagander)

  This simplifies setting up a standby server.

- Allow `pg_receivexlog` and `pg_basebackup --xlog-method` to handle streaming timeline switches (Heikki Linnakangas)

- Add `wal_receiver_timeout` parameter to control the WAL receiver's timeout (Amit Kapila)

  This allows more rapid detection of connection failure.

- Change the WAL record format to allow splitting the record header across pages (Heikki Linnakangas)

  The new format is slightly more compact, and is more efficient to write.

E.84.3.3. Queries

- Implement SQL-standard `LATERAL` option for `FROM`-clause subqueries and function calls (Tom Lane)

  This feature allows subqueries and functions in `FROM` to reference columns from other tables in the `FROM` clause. The `LATERAL` keyword is optional for functions.

- Add support for piping `COPY` and `psql \copy` data to/from an external program (Etsuro Fujita)

- Allow a multirow `VALUES` clause in a rule to reference `OLD/NEW` (Tom Lane)

E.84.3.4. Object Manipulation

- Add support for event triggers (Dimitri Fontaine, Robert Haas, Álvaro Herrera)

  This allows server-side functions written in event-enabled languages to be called when DDL commands are run.
• Allow foreign data wrappers to support writes (inserts/updates/deletes) on foreign tables (KaiGai Kohei)

• Add CREATE SCHEMA ... IF NOT EXISTS clause (Fabrízio de Royes Mello)

• Make REASSIGN OWNED also change ownership of shared objects (Álvaro Herrera)

• Make CREATE AGGREGATE complain if the given initial value string is not valid input for the transition datatype (Tom Lane)

• Suppress CREATE TABLE's messages about implicit index and sequence creation (Robert Haas)
  These messages now appear at DEBUG1 verbosity, so that they will not be shown by default.

• Allow DROP TABLE IF EXISTS to succeed when a non-existent schema is specified in the table name (Bruce Momjian)
  Previously, it threw an error if the schema did not exist.

• Provide clients with constraint violation details as separate fields (Pavel Stehule)
  This allows clients to retrieve table, column, data type, or constraint name error details. Previously such information had to be extracted from error strings. Client library support is required to access these fields.

**E.84.3.4.1. ALTER**

• Support IF NOT EXISTS option in ALTER TYPE ... ADD VALUE (Andrew Dunstan)
  This is useful for conditionally adding values to enumerated types.

• Add ALTER ROLE ALL SET to establish settings for all users (Peter Eisentraut)
  This allows settings to apply to all users in all databases. ALTER DATABASE SET already allowed addition of settings for all users in a single database. postgresql.conf has a similar effect.

• Add support for ALTER RULE ... RENAME (Ali Dar)

**E.84.3.4.2. VIEWS**

• Add materialized views (Kevin Grittner)
  Unlike ordinary views, where the base tables are read on every access, materialized views create physical tables at creation or refresh time. Access to the materialized view then reads from its physical table. There is not yet any facility for incrementally refreshing materialized views or auto-accessing them via base table access.

• Make simple views auto-updatable (Dean Rasheed)
  Simple views that reference some or all columns from a single base table are now updatable by default. More complex views can be made updatable using INSTEAD OF triggers or INSTEAD rules.

• Add CREATE RECURSIVE VIEW syntax (Peter Eisentraut)
  Internally this is translated into CREATE VIEW ... WITH RECURSIVE ....

• Improve view/rule printing code to handle cases where referenced tables are renamed, or columns are renamed, added, or dropped (Tom Lane)
  Table and column renamings can produce cases where, if we merely substitute the new name into the original text of a rule or view, the result is ambiguous. This change fixes the rule-dumping code to insert manufactured table and column aliases when needed to preserve the original semantics.
E.84.3.5. Data Types

- Increase the maximum size of large objects from 2GB to 4TB (Nozomi Anzai, Yugo Nagata)
  This change includes adding 64-bit-capable large object access functions, both in the server and
  in libpq.
- Allow text timezone designations, e.g. “America/Chicago”, in the “T” field of ISO-format
  timestamp input (Bruce Momjian)

E.84.3.5.1. JSON

- Add operators and functions to extract elements from JSON values (Andrew Dunstan)
- Allow JSON values to be converted into records (Andrew Dunstan)
- Add functions to convert scalars, records, and hstore values to JSON (Andrew Dunstan)

E.84.3.6. Functions

- Add array_remove() and array_replace() functions (Marco Nenciarini, Gabriele Bartolini)
- Allow concat() and format() to properly expand VARIADIC-labeled arguments (Pavel Stehule)
- Improve format() to provide field width and left/right alignment options (Pavel Stehule)
- Make to_char(), to_date(), and to_timestamp() handle negative (BC) century values
  properly (Bruce Momjian)
  Previously the behavior was either wrong or inconsistent with positive/AD handling, e.g. with the
  format mask “IYYYY-IW-DY”.
- Make to_date() and to_timestamp() return proper results when mixing ISO and Gregorian
  week/day designations (Bruce Momjian)
- Cause pg_get_viewdef() to start a new line by default after each SELECT target list entry and
  FROM entry (Marko Tiikkaja)
  This reduces line length in view printing, for instance in pg_dump output.
- Fix map_sql_value_to_xml_value() to print values of domain types the same way their
  base type would be printed (Pavel Stehule)
  There are special formatting rules for certain built-in types such as boolean; these rules now also
  apply to domains over these types.

E.84.3.7. Server-Side Languages

E.84.3.7.1. PL/pgSQL Server-Side Language

- Allow PL/pgSQL to use RETURN with a composite-type expression (Asif Rehman)
  Previously, in a function returning a composite type, RETURN could only reference a variable of
  that type.
- Allow PL/pgSQL to access constraint violation details as separate fields (Pavel Stehule)
- Allow PL/pgSQL to access the number of rows processed by COPY (Pavel Stehule)
A COPY executed in a PL/pgSQL function now updates the value retrieved by GET DIAGNOSTICS
\[ x = \text{ROW\_COUNT}. \]

- Allow unreserved keywords to be used as identifiers everywhere in PL/pgSQL (Tom Lane)

In certain places in the PL/pgSQL grammar, keywords had to be quoted to be used as identifiers, even if they were nominally unreserved.

**E.84.3.7.2. PL/Python Server-Side Language**

- Add PL/Python result object string handler (Peter Eisentraut)
  
  This allows `plpy.debug(rv)` to output something reasonable.

- Make PL/Python convert OID values to a proper Python numeric type (Peter Eisentraut)

- Handle SPI errors raised explicitly (with PL/Python’s `RAISE`) the same as internal SPI errors (Oskari Saarenmaa and Jan Urbanski)

**E.84.3.8. Server Programming Interface (SPI)**

- Prevent leakage of SPI tuple tables during subtransaction abort (Tom Lane)

  At the end of any failed subtransaction, the core SPI code now releases any SPI tuple tables that were created during that subtransaction. This avoids the need for SPI-using code to keep track of such tuple tables and release them manually in error-recovery code. Failure to do so caused a number of transaction-lifespan memory leakage issues in PL/pgSQL and perhaps other SPI clients. `SPI_freetuptable()` now protects itself against multiple freeing requests, so any existing code that did take care to clean up shouldn’t be broken by this change.

- Allow SPI functions to access the number of rows processed by COPY (Pavel Stehule)

**E.84.3.9. Client Applications**

- Add command-line utility `pg_isready` to check if the server is ready to accept connections (Phil Sorber)

- Support multiple `--table` arguments for `pg_restore`, `clusterdb`, `reindexdb`, and `vacuumdb` (Josh Kupershmidt)

  This is similar to the way `pg_dump`’s `--table` option works.

- Add `--dbname` option to `pg_dumpall`, `pg_basebackup`, and `pg_receivexlog` to allow specifying a connection string (Amit Kapila)

- Add libpq function `PQconninfo()` to return connection information (Zoltán Bőszörményi, Magnus Hagander)

**E.84.3.9.1. psql**

- Adjust function cost settings so `psql` tab completion and pattern searching are more efficient (Tom Lane)

- Improve `psql`’s tab completion coverage (Jeff Janes, Dean Rasheed, Peter Eisentraut, Magnus Hagander)

- Allow the `psql` `--single-transaction` mode to work when reading from standard input (Fabien Coelho, Robert Haas)

  Previously this option only worked when reading from a file.
• Remove psql warning when connecting to an older server (Peter Eisentraut)

A warning is still issued when connecting to a server of a newer major version than psql's.

E.84.3.9.1.1. Backslash Commands

• Add psql command \watch to repeatedly execute a SQL command (Will Leinweber)
• Add psql command \gset to store query results in psql variables (Pavel Stehule)
• Add SSL information to psql's \conninfo command (Alastair Turner)
• Add “Security” column to psql's \df+ output (Jon Erdman)
• Allow psql command \l to accept a database name pattern (Peter Eisentraut)
• In psql, do not allow \connect to use defaults if there is no active connection (Bruce Momjian)
  This might be the case if the server had crashed.
• Properly reset state after failure of a SQL command executed with psql's \g file (Tom Lane)
  Previously, the output from subsequent SQL commands would unexpectedly continue to go to the same file.

E.84.3.9.1.2. Output

• Add a latex-longtable output format to psql (Bruce Momjian)
  This format allows tables to span multiple pages.
• Add a border=3 output mode to the psql latex format (Bruce Momjian)
• In psql's tuples-only and expanded output modes, no longer emit “(No rows)” for zero rows (Peter Eisentraut)
• In psql's unaligned, expanded output mode, no longer print an empty line for zero rows (Peter Eisentraut)

E.84.3.9.2. pg_dump

• Add pg_dump --jobs option to dump tables in parallel (Joachim Wieland)
• Make pg_dump output functions in a more predictable order (Joel Jacobson)
• Fix tar files emitted by pg_dump to be POSIX conformant (Brian Weaver, Tom Lane)
• Add --dbname option to pg_dump, for consistency with other client commands (Heikki Linnakangas)
  The database name could already be supplied last without a flag.

E.84.3.9.3. initdb

• Make initdb fsync the newly created data directory (Jeff Davis)
  This insures data integrity in event of a system crash shortly after initdb. This can be disabled by using --nosync.
• Add initdb --sync-only option to sync the data directory to durable storage (Bruce Momjian)
  This is used by pg_upgrade.
- Make initdb issue a warning about placing the data directory at the top of a file system mount point (Bruce Momjian)

### E.84.3.10. Source Code

- Add infrastructure to allow plug-in background worker processes (Álvaro Herrera)
- Create a centralized timeout API (Zoltán Bőszörményi)
- Create libpgcommon and move `pg_malloc()` and other functions there (Álvaro Herrera, Andres Freund)
  
  This allows libpqport to be used solely for portability-related code.
- Add support for list links embedded in larger structs (Andres Freund)
- Use `SA_RESTART` for all signals, including `SIGALRM` (Tom Lane)
- Ensure that the correct text domain is used when translating `errcontext()` messages (Heikki Linnakangas)
- Standardize naming of client-side memory allocation functions (Tom Lane)
- Provide support for “static assertions” that will fail at compile time if some compile-time-constant condition is not met (Andres Freund, Tom Lane)
- Support `Assert()` in client-side code (Andrew Dunstan)
- Add decoration to inform the C compiler that some `ereport()` and `elog()` calls do not return (Peter Eisentraut, Andres Freund, Tom Lane, Heikki Linnakangas)
- Allow options to be passed to the regression test output comparison utility via `PG_REGRESS_DIFF_OPTS` (Peter Eisentraut)
- Add isolation tests for `CREATE INDEX CONCURRENTLY` (Abhijit Menon-Sen)
- Remove typedefs for `int2/int4` as they are better represented as `int16/int32` (Peter Eisentraut)
- Fix `install-strip` on Mac OS X (Peter Eisentraut)
- Remove configure flag `--disable-shared`, as it is no longer supported (Bruce Momjian)
- Rewrite `pgindent` in Perl (Andrew Dunstan)
- Provide Emacs macro to set Perl formatting to match PostgreSQL’s `perltidy` settings (Peter Eisentraut)
- Run tool to check the keyword list whenever the backend grammar is changed (Tom Lane)
- Change the way `UESCAPE` is lexed, to significantly reduce the size of the lexer tables (Heikki Linnakangas)
- Centralize flex and bison make rules (Peter Eisentraut)

  This is useful for `pgxs` authors.
- Change many internal backend functions to return object `OIDs` rather than void (Dimitri Fontaine)

  This is useful for event triggers.
- Invent pre-commit/pre-prepare/pre-subcommit events for transaction callbacks (Tom Lane)
Loadable modules that use transaction callbacks might need modification to handle these new event types.

- Add function `pg_identify_object()` to produce a machine-readable description of a database object (Álvaro Herrera)
- Add post-ALTER-object server hooks (KaiGai Kohei)
- Implement a generic binary heap and use it for Merge-Append operations (Abhijit Menon-Sen)
- Provide a tool to help detect timezone abbreviation changes when updating the src/timezone/data files (Tom Lane)
- Add pkg-config support for libpq and ecpg libraries (Peter Eisentraut)
- Remove src/tools/backend, now that the content is on the PostgreSQL wiki (Bruce Momjian)
- Split out WAL reading as an independent facility (Heikki Linnakangas, Andres Freund)
- Use a 64-bit integer to represent WAL positions (`XLogRecPtr`) instead of two 32-bit integers (Heikki Linnakangas)

Generally, tools that need to read the WAL format will need to be adjusted.

- Allow PL/Python to support platform-specific include directories (Peter Eisentraut)
- Allow PL/Python on OS X to build against custom versions of Python (Peter Eisentraut)

E.84.3.11. Additional Modules

- Add a Postgres foreign data wrapper contrib module to allow access to other Postgres servers (Shigeru Hanada)

  This foreign data wrapper supports writes.

- Add `pg_xlogdump` contrib program (Andres Freund)
- Add support for indexing of regular-expression searches in `pg_trgm` (Alexander Korotkov)
- Improve `pg_trgm`'s handling of multibyte characters (Tom Lane)

  On a platform that does not have the wcstombs() or tolower() library functions, this could result in an incompatible change in the contents of `pg_trgm` indexes for non-ASCII data. In such cases, `REINDEX` those indexes to ensure correct search results.

- Add a `pgstattuple` function to report the size of the pending-insertions list of a GIN index (Fujii Masao)
- Make `oid2name`, `pgbench`, and `vacuumlo` set `fallback_application_name` (Amit Kapila)
- Improve output of `pg_test_timing` (Bruce Momjian)
- Improve output of `pg_test_fsync` (Peter Geoghegan)
- Create a dedicated foreign data wrapper, with its own option validator function, for `dblink` (Shigeru Hanada)

When using this FDW to define the target of a `dblink` connection, instead of using a hard-wired list of connection options, the underlying libpq library is consulted to see what connection options it supports.
E.84.3.11.1. pg_upgrade

- Allow pg_upgrade to do dumps and restores in parallel (Bruce Momjian, Andrew Dunstan)
  This allows parallel schema dump/restore of databases, as well as parallel copy/link of data files per tablespace. Use the --jobs option to specify the level of parallelism.
- Make pg_upgrade create Unix-domain sockets in the current directory (Bruce Momjian, Tom Lane)
  This reduces the possibility that someone will accidentally connect during the upgrade.
- Make pg_upgrade --check mode properly detect the location of non-default socket directories (Bruce Momjian, Tom Lane)
- Improve performance of pg_upgrade for databases with many tables (Bruce Momjian)
- Improve pg_upgrade's logs by showing executed commands (Álvaro Herrera)
- Improve pg_upgrade's status display during copy/link (Bruce Momjian)

E.84.3.11.2. pgbench

- Add --foreign-keys option to pgbench (Jeff Janes)
  This adds foreign key constraints to the standard tables created by pgbench, for use in foreign key performance testing.
- Allow pgbench to aggregate performance statistics and produce output every --aggregate-interval seconds (Tomas Vondra)
- Add pgbench --sampling-rate option to control the percentage of transactions logged (Tomas Vondra)
- Reduce and improve the status message output of pgbench's initialization mode (Robert Haas, Peter Eisentraut)
- Add pgbench -q mode to print one output line every five seconds (Tomas Vondra)
- Output pgbench elapsed and estimated remaining time during initialization (Tomas Vondra)
- Allow pgbench to use much larger scale factors, by changing relevant columns from integer to bigint when the requested scale factor exceeds 20000 (Greg Smith)

E.84.3.12. Documentation

- Allow EPUB-format documentation to be created (Peter Eisentraut)
- Update FreeBSD kernel configuration documentation (Brad Davis)
- Improve WINDOW function documentation (Bruce Momjian, Florian Pflug)
- Add instructions for setting up the documentation tool chain on macOS (Peter Eisentraut)
- Improve commit_delay documentation (Peter Geoghegan)

E.85. Release 9.2.24

Release date: 2017-11-09

This release contains a variety of fixes from 9.2.23. For information about new features in the 9.2 major release, see Section E.109.
This is expected to be the last PostgreSQL release in the 9.2.X series. Users are encouraged to update to a newer release branch soon.

### E.85.1. Migration to Version 9.2.24

A dump/restore is not required for those running 9.2.X. However, if you are upgrading from a version earlier than 9.2.22, see Section E.87.

### E.85.2. Changes

- **Fix sample server-start scripts to become \$PGUSER before opening \$PGLOG** (Noah Misch)

  Previously, the postmaster log file was opened while still running as root. The database owner could therefore mount an attack against another system user by making \$PGLOG be a symbolic link to some other file, which would then become corrupted by appending log messages.

  By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing \$PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

- **Properly reject attempts to convert infinite float values to type numeric** (Tom Lane, KaiGai Kohei)

  Previously the behavior was platform-dependent.

- **Fix corner-case crashes when columns have been added to the end of a view** (Tom Lane)

- **Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes** (Tom Lane)

  Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

- **Correctly detect hashability of range data types** (Tom Lane)

  The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range's subtype has hash support. This does not affect any of the built-in range types, since they're all hashable anyway.

- **Fix low-probability loss of NOTIFY messages due to XID wraparound** (Marko Tiikkaja, Tom Lane)

  If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.

- **Prevent low-probability crash in processing of nested trigger firings** (Tom Lane)

- **Correctly restore the umask setting when file creation fails in COPY or lo_export()** (Peter Eisentraut)

- **Give a better error message for duplicate column names in ANALYZE** (Nathan Bossart)

- **Fix libpq to not require user's home directory to exist** (Tom Lane)

  In v10, failure to find the home directory while trying to read ~/.pgpass was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading ~/.pg_service.conf, though this was less obvious since that file is not sought unless a service name is specified.
• Fix libpq to guard against integer overflow in the row count of a PGresult (Michael Paquier)

• Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

This fixes various issues; the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no posixrules file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.

• Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

E.86. Release 9.2.23

Release date: 2017-08-31

This release contains a small number of fixes from 9.2.22. For information about new features in the 9.2 major release, see Section E.109.

The PostgreSQL community will stop releasing updates for the 9.2.X release series in September 2017. Users are encouraged to update to a newer release branch soon.

E.86.1. Migration to Version 9.2.23

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.22, see Section E.87.

E.86.2. Changes

• Show foreign tables in information_schema.table_privileges view (Peter Eisentraut)

All other relevant information_schema views include foreign tables, but this one ignored them.

Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:

```sql
SET search_path TO information_schema;
CREATE OR REPLACE VIEW table_privileges AS
  SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
         CAST(grantee.rolname AS sql_identifier) AS grantee,
         CAST(current_database() AS sql_identifier) AS table_catalog,
         CAST(nc.nspname AS sql_identifier) AS table_schema,
         CAST(c.relname AS sql_identifier) AS table_name,
         CAST(c.prtype AS character_data) AS privilege_type,
         CAST(
          CASE WHEN
            -- object owner always has grant options
            pg_has_role(grantee.oid, c.relowner, 'USAGE')
          OR c.grantable
          THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
         CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE 'NO' END AS yes_or_no) AS with_hierarchy
```
FROM (  
    SELECT oid, relname, relnamespace, relkind, relowner,  
    (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class  
) AS c (oid, relname, relnamespace, relkind, relowner,  
grantor, grantee, prtype, grantable),  
pg_namespace nc,  
pag_authid u_grantor,  

    SELECT oid, rolname FROM pag_authid  
UNION ALL  
SELECT 0::oid, 'PUBLIC'  
) AS grantee (oid, rolname)

WHERE c.relnamespace = nc.oid  
AND c.relkind IN ('r', 'v', 'f')  
AND c.grantee = grantee.oid  
AND c.grantor = u_grantor.oid  
AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE', 'DELETE', 'REFERENCES', 'TRIGGER')  
AND (pg_has_role(u_grantor.oid, 'USAGE')  
OR pg_has_role(grantee.oid, 'USAGE')  
OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

• Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)

This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about "cannot drop active portal".

• Remove assertion that could trigger during a fatal exit (Tom Lane)

• Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant values that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.

• Change ecpg's parser to allow RETURNING clauses without attached C variables (Michael Meskes)

This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

• Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that _USE_32BIT_TIME_T is now longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

E.87. Release 9.2.22

Release date: 2017-08-10
This release contains a variety of fixes from 9.2.21. For information about new features in the 9.2 major release, see Section E.109.

The PostgreSQL community will stop releasing updates for the 9.2.X release series in September 2017. Users are encouraged to update to a newer release branch soon.

**E.87.1. Migration to Version 9.2.22**

A dump/restore is not required for those running 9.2.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.2.20, see Section E.89.

**E.87.2. Changes**

- Further restrict visibility of `pg_user_mappings.umoptions`, to protect passwords stored as user mapping options (Noah Misch)

The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have `USAGE` permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since `information_schema.user_mapping_options` does not show the options in such cases, `pg_user_mappings` should not either. (CVE-2017-7547)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this change in an existing database, you will need to do the following:

1. Restart the postmaster after adding `allow_system_table_mods = true` to `postgresql.conf`. (In versions supporting `ALTER SYSTEM`, you can use that to make the configuration change, but you'll still need a restart.)

2. In each database of the cluster, run the following commands as superuser:

```sql
SET search_path = pg_catalog;
CREATE OR REPLACE VIEW pg_user_mappings AS
SELECT
    U.oid       AS umid,
    S.oid       AS srvid,
    S.srvname   AS srvname,
    U.umuser    AS umuser,
    CASE WHEN U.umuser = 0 THEN
        'public'
    ELSE
        A.rolname
    END AS usename,
    CASE WHEN (U.umuser <> 0 AND A.rolname = current_user
        AND (pg_has_role(S.srvowner, 'USAGE')
            OR has_server_privilege(S.oid,
            'USAGE')))
        OR (U.umuser = 0 AND
        pg_has_role(S.srvowner, 'USAGE'))
        OR (SELECT rolsuper FROM pg_authid WHERE
        rolname = current_user)
        THEN U.umoptions
    ELSE NULL END AS umoptions
FROM pg_user_mapping U
```
LEFT JOIN pg_authid A ON (A.oid = U.umuser) JOIN pg_foreign_server S ON (U.umserver = S.oid);

3. Do not forget to include the template0 and template1 databases, or the vulnerability will still exist in databases you create later. To fix template0, you'll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

```
ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;
```

and then after fixing template0, undo that with

```
ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;
```

In prior versions, instead use

```
UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';
```

4. Finally, remove the allow_system_table_mods configuration setting, and again restart the postmaster.

• Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)

  libpq ignores empty password specifications, and does not transmit them to the server. So, if a user's password has been set to the empty string, it's impossible to log in with that password via psql or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, md5, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

• On Windows, retry process creation if we fail to reserve the address range for our shared memory in the new process (Tom Lane, Amit Kapila)

  This is expected to fix infrequent child-process-launch failures that are probably due to interference from antivirus products.

• Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas Munro, Tom Lane)

• Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael Paquier)

• Prevent sending SSL session tickets to clients (Tom Lane)

  This fix prevents reconnection failures with ticket-aware client-side SSL code.

• Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)

• Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and immediate restart (Tom Lane)

  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively ignored.

• Ensure that the statistics collector's receive buffer size is at least 100KB (Tom Lane)
This reduces the risk of dropped statistics data on older platforms whose default receive buffer size is less than that.

- Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes an XLOG_SWITCH WAL record (Andres Freund)

- Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)

- Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom Lane)

- Fix cases where an INSERT or UPDATE assigns to more than one element of a column that is of domain-over-array type (Tom Lane)

- Move autogenerated array types out of the way during ALTER ... RENAME (Vik Fearing)

  Previously, we would rename a conflicting autogenerated array type out of the way during CREATE; this fix extends that behavior to renaming operations.

- Ensure that ALTER USER ... SET accepts all the syntax variants that ALTER ROLE ... SET does (Peter Eisentraut)

- Properly update dependency info when changing a datatype I/O function's argument or return type from opaque to the correct type (Heikki Linnakangas)

  CREATE TYPE updates I/O functions declared in this long-obsolete style, but it forgot to record a dependency on the type, allowing a subsequent DROP TYPE to leave broken function definitions behind.

- Reduce memory usage when ANALYZE processes a tsvector column (Heikki Linnakangas)

- Fix unnecessary precision loss and sloppy rounding when multiplying or dividing money values by integers or floats (Tom Lane)

- Tighten checks for whitespace in functions that parse identifiers, such as regprocedurein() (Tom Lane)

  Depending on the prevailing locale, these functions could misinterpret fragments of multibyte characters as whitespace.

- Use relevant #define symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)

  This avoids portability problems, typically manifesting as a “handshake” mismatch during library load, when working with recent Perl versions.

- In psql, fix failure when COPY FROM STDIN is ended with a keyboard EOF signal and then another COPY FROM STDIN is attempted (Thomas Munro)

  This misbehavior was observed on BSD-derived platforms (including macOS), but not on most others.

- Fix pg_dump to not emit invalid SQL for an empty operator class (Daniel Gustafsson)

- Fix pg_dump output to stdout on Windows (Kuntal Ghosh)

  A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.

- Fix pg_get_ruledef() to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)
In some corner cases, pg_dump relies on `pg_get_ruledef()` to dump views, so that this error could result in dump/reload failures.

- Fix dumping of function expressions in the `FROM` clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)
- Fix `pg_basebackup` output to stdout on Windows (Haribabu Kommi)
  
  A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.
- Fix `pg_upgrade` to ensure that the ending WAL record does not have `wal_level = minimum` (Bruce Momjian)
  
  This condition could prevent upgraded standby servers from reconnecting.
- Always use `-fPIC`, not `-fpic`, when building shared libraries with gcc (Tom Lane)
  
  This supports larger extension libraries on platforms where it makes a difference.
- Fix unescaped-braces issue in our build scripts for Microsoft MSVC, to avoid a warning or error from recent Perl versions (Andrew Dunstan)
- In MSVC builds, handle the case where the OpenSSL library is not within a VC subdirectory (Andrew Dunstan)
- In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)
  
  This fixes a former need to move things around in standard Windows installations of libxml2.
- In MSVC builds, recognize a Tcl library that is named `tcl86.lib` (Noah Misch)

**E.88. Release 9.2.21**

**Release date:** 2017-05-11

This release contains a variety of fixes from 9.2.20. For information about new features in the 9.2 major release, see Section E.109.

The PostgreSQL community will stop releasing updates for the 9.2.X release series in September 2017. Users are encouraged to update to a newer release branch soon.

**E.88.1. Migration to Version 9.2.21**

A dump/restore is not required for those running 9.2.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.2.20, see Section E.89.

**E.88.2. Changes**

- Restrict visibility of `pg_user_mappings.umoptions`, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)
  
  The previous coding allowed the owner of a foreign server object, or anyone he has granted server `USAGE` permission to, to see the options for all user mappings associated with that server. This
might well include passwords for other users. Adjust the view definition to match the behavior of
information_schema.user_mapping_options, namely that these options are visible to
the user being mapped, or if the mapping is for PUBLIC and the current user is the server owner,
or if the current user is a superuser. (CVE-2017-7486)

By itself, this patch will only fix the behavior in newly initdb'd databases. If you wish to apply this
change in an existing database, follow the corrected procedure shown in the changelog entry for
CVE-2017-7547, in Section E.87.

- Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

Some selectivity estimation functions in the planner will apply user-defined operators to values
obtained from pg_statistic, such as most common values and histogram entries. This occurs
before table permissions are checked, so a nefarious user could exploit the behavior to obtain these
values for table columns he does not have permission to read. To fix, fall back to a default estimate
if the operator's implementation function is not certified leak-proof and the calling user does not
have permission to read the table column whose statistics are needed. At least one of these criteria
is satisfied in most cases in practice. (CVE-2017-7484)

- Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)

This could result in an unlogged index being set to an invalid state after a crash and restart. Such a
problem would persist until the index was dropped and rebuilt.

- Fix incorrect reconstruction of pg_subtrans entries when a standby server replays a prepared
  but uncommitted two-phase transaction (Tom Lane)

In most cases this turned out to have no visible ill effects, but in corner cases it could result in
circular references in pg_subtrans, potentially causing infinite loops in queries that examine
rows modified by the two-phase transaction.

- Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL.(Julien
  Rouhaud, Tom Lane)

Due to lack of a cache flush step between commands in an extension script file, non-utility queries
might not see the effects of an immediately preceding catalog change, such as ALTER TABLE ... 
RENAME.

- Skip tablespace privilege checks when ALTER TABLE ... ALTER COLUMN TYPE rebuilds
  an existing index (Noah Misch)

The command failed if the calling user did not currently have CREATE privilege for the tablespace
containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be
rebuilt where it is.

- Fix ALTER TABLE ... VALIDATE CONSTRAINT to not recurse to child tables when the
  constraint is marked NO INHERIT (Amit Langote)

This fix prevents unwanted “constraint does not exist” failures when no matching constraint is
present in the child tables.

- Fix VACUUM to account properly for pages that could not be scanned due to conflicting page pins
  (Andrew Gierth)

This tended to lead to underestimation of the number of tuples in the table. In the worst case of a
small heavily-contended table, VACUUM could incorrectly report that the table contained no tuples,
leading to very bad planning choices.

- Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests
  (Tom Lane, Thomas Munro)
• Fix `cursor_to_xml()` to produce valid output with `tableforest = false` (Thomas Munro, Peter Eisentraut)

Previously it failed to produce a wrapping `<table>` element.

• Improve performance of `pg_timezone_names` view (Tom Lane, David Rowley)

• Fix sloppy handling of corner-case errors from `lseek()` and `close()` (Tom Lane)

Neither of these system calls are likely to fail in typical situations, but if they did, `fd.c` could get quite confused.

• Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)

This could result in attempting to write to the event log when that isn't accessible, so that no logging happens at all.

• Fix `ecpg` to support `COMMIT PREPARED` and `ROLLBACK PREPARED` (Masahiko Sawada)

• Fix a double-free error when processing dollar-quoted string literals in `ecpg` (Michael Meskes)

• In `pg_dump`, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)

In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.

• Avoid emitting an invalid list file in `pg_restore -l` when SQL object names contain newlines (Tom Lane)

Replace newlines by spaces, which is sufficient to make the output valid for `pg_restore -L`'s purposes.

• Fix `pg_upgrade` to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)

Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.

• Improve error handling in `contrib/adminpack's pg_file_write()` function (Noah Misch)

Notably, it failed to detect errors reported by `fclose()`.

• In `contrib/dblink`, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

• Support OpenSSL 1.1.0 (Heikki Linnakangas, Andreas Karlsson, Tom Lane)

This is a back-patch of work previously done in newer branches; it's needed since many platforms are adopting newer OpenSSL versions.

• Support Tcl 8.6 in MSVC builds (Álvaro Herrera)

• Sync our copy of the timezone library with IANA release `tzcode2017b` (Tom Lane)

This fixes a bug affecting some DST transitions in January 2038.

• Update time zone data files to `t泽data release 2017b` for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric
abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

- Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds (David Rowley)

The Microsoft MSVC build scripts neglected to install the posixrules file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the US/Eastern zone. If you don't want that, remove the posixrules file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.

E.89. Release 9.2.20

Release date: 2017-02-09

This release contains a variety of fixes from 9.2.19. For information about new features in the 9.2 major release, see Section E.109.

E.89.1. Migration to Version 9.2.20

A dump/restore is not required for those running 9.2.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.89.2. Changes

- Fix a race condition that could cause indexes built with CREATE INDEX CONCURRENTLY to be corrupt (Pavan Deolasee, Tom Lane)

If CREATE INDEX CONCURRENTLY was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the CREATE INDEX command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

- Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)

Previously, this was skipped when wal_level = minimal, but actually it's necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

- Fix WAL page header validation when re-reading segments (Takayuki Tsunakawa, Amit Kapila)

In corner cases, a spurious “out-of-sequence TLI” error could be reported during recovery.

- If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)
• Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

• Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)

This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

• Fix check for when an extension member object can be dropped (Tom Lane)

Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

• Make sure ALTER TABLE preserves index tablespace assignments when rebuilding indexes (Tom Lane, Michael Paquier)

Previously, non-default settings of default_tablespace could result in broken indexes.

• Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)

This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.

• Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via ALTER TABLE ... INHERIT (Amit Langote)

The OID column should be treated the same as regular user columns in this case, but it wasn’t, leading to odd behavior in later inheritance changes.

• Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

• Ensure that column typmods are determined accurately for multi-row VALUES constructs (Tom Lane)

This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a varchar value) but later values don't share the same limit.

• Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (U&'...') or Unicode identifier (U&"...").

• Ensure that a purely negative text search query, such as !foo, matches empty tsvectors (Tom Dunstan)

Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

• Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

• Fix performance problems in ts_rewrite() (Tom Lane)

• Fix ts_rewrite()'s handling of nested NOT operators (Tom Lane)
• Fix `array_fill()` to handle empty arrays properly (Tom Lane)

• Fix one-byte buffer overrun in `quote_literal_cstr()` (Heikki Linnakangas)

  The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.

• Prevent multiple calls of `pg_start_backup()` and `pg_stop_backup()` from running concurrently (Michael Paquier)

  This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

• Avoid discarding `interval`-to-`interval` casts that aren't really no-ops (Tom Lane)

  In some cases, a cast that should result in zeroing out low-order `interval` fields was mistakenly deemed to be a no-op and discarded. An example is that casting from `INTERVAL` `MONTH` to `INTERVAL` `YEAR` failed to clear the months field.

• Fix `pg_dump` to dump user-defined casts and transforms that use built-in functions (Stephen Frost)

• Fix possible `pg_basebackup` failure on standby server when including WAL files (Amit Kapila, Robert Haas)

• Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)

  This avoids failures if the objects are used after a Python garbage collection cycle has occurred.

• Fix PL/Tcl to support triggers on tables that have `.tupno` as a column name (Tom Lane)

  This matches the (previously undocumented) behavior of PL/Tcl's `spi_exec` and `spi_execp` commands, namely that a magic `.tupno` column is inserted only if there isn't a real column named that.

• Allow DOS-style line endings in `~/.pgpass` files, even on Unix (Vik Fearing)

  This change simplifies use of the same password file across Unix and Windows machines.

• Fix one-byte buffer overrun if `ecpg` is given a file name that ends with a dot (Takayuki Tsunakawa)

• Fix `psql`'s tab completion for `ALTER DEFAULT PRIVILEGES` (Gilles Darold, Stephen Frost)

• In `psql`, treat an empty or all-blank setting of the `PAGER` environment variable as meaning “no pager” (Tom Lane)

  Previously, such a setting caused output intended for the pager to vanish entirely.

• Improve contrib/dblink's reporting of low-level `libpq` errors, such as out-of-memory (Joe Conway)

• On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)

• Sync our copy of the timezone library with IANA release `tzcode2016j` (Tom Lane)

  This fixes various issues, most notably that timezone data installation failed if the target directory didn't support hard links.

• Update time zone data files to `tzdata` release `2016j` for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarctica/Casery. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.
E.90. Release 9.2.19

Release date: 2016-10-27

This release contains a variety of fixes from 9.2.18. For information about new features in the 9.2 major release, see Section E.109.

E.90.1. Migration to Version 9.2.19

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.90.2. Changes

- Fix EvalPlanQual rechecks involving CTE scans (Tom Lane)

  The recheck would always see the CTE as returning no rows, typically leading to failure to update rows that were recently updated.

- Fix improper repetition of previous results from hashed aggregation in a subquery (Andrew Gierth)

  The test to see if we can reuse a previously-computed hash table of the aggregate state values neglected the possibility of an outer query reference appearing in an aggregate argument expression. A change in the value of such a reference should lead to recalculating the hash table, but did not.

- Fix EXPLAIN to emit valid XML when track_io_timing is on (Markus Winand)

  Previously the XML output-format option produced syntactically invalid tags such as <I/O-Read-Time>. That is now rendered as <I-O-Read-Time>.

- Suppress printing of zeroes for unmeasured times in EXPLAIN (Maksim Milyutin)

  Certain option combinations resulted in printing zero values for times that actually aren’t ever measured in that combination. Our general policy in EXPLAIN is not to print such fields at all, so do that consistently in all cases.

- Fix timeout length when VACUUM is waiting for exclusive table lock so that it can truncate the table (Simon Riggs)

  The timeout was meant to be 50 milliseconds, but it was actually only 50 microseconds, causing VACUUM to give up on truncation much more easily than intended. Set it to the intended value.

- Fix bugs in merging inherited CHECK constraints while creating or altering a table (Tom Lane, Amit Langote)

  Allow identical CHECK constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a NOT VALID constraint on the child. Likewise, prevent merging of a NO INHERIT child constraint with an inherited constraint.

- Remove artificial restrictions on the values accepted by numeric_in() and numeric_recv() (Tom Lane)

  We allow numeric values up to the limit of the storage format (more than $1e100000$), so it seems fairly pointless that numeric_in() rejected scientific-notation exponents above 1000. Likewise, it was silly for numeric_recv() to reject more than 1000 digits in an input value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)
• Fix file descriptor leakage when truncating a temporary relation of more than 1GB (Andres Freund)

• Disallow starting a standalone backend with standby_mode turned on (Michael Paquier)

This can't do anything useful, since there will be no WAL receiver process to fetch more WAL data; and it could result in misbehavior in code that wasn't designed with this situation in mind.

• Don't try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)

This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)

The reported problem involved leaking an error report during PQreset(), but there might be related cases.

• Make ecpg's --help and --version options work consistently with our other executables (Haribabu Kommi)

• In pg_dump, never dump range constructor functions (Tom Lane)

This oversight led to pg_upgrade failures with extensions containing range types, due to duplicate creation of the constructor functions.

• Fix contrib/intarray/bench/bench.pl to print the results of the EXPLAIN it does when given the -e option (Daniel Gustafsson)

• Update Windows time zone mapping to recognize some time zone names added in recent Windows versions (Michael Paquier)

• Prevent failure of obsolete dynamic time zone abbreviations (Tom Lane)

If a dynamic time zone abbreviation does not match any entry in the referenced time zone, treat it as equivalent to the time zone name. This avoids unexpected failures when IANA removes abbreviations from their time zone database, as they did in tzdata release 2016f and seem likely to do again in the future. The consequences were not limited to not recognizing the individual abbreviation; any mismatch caused the pg_timezone_abbrevs view to fail altogether.

• Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.91. Release 9.2.18

Release date: 2016-08-11

This release contains a variety of fixes from 9.2.17. For information about new features in the 9.2 major release, see Section E.109.
E.91.1. Migration to Version 9.2.18

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.91.2. Changes

• Fix possible mis-evaluation of nested CASE-WHEN expressions (Heikki Linnakangas, Michael Paquier, Tom Lane)

A CASE expression appearing within the test value subexpression of another CASE could become confused about whether its own test value was null or not. Also, inlining of a SQL function implementing the equality operator used by a CASE expression could result in passing the wrong test value to functions called within a CASE expression in the SQL function's body. If the test values were of different data types, a crash might result; moreover such situations could be abused to allow disclosure of portions of server memory. (CVE-2016-5423)

• Fix client programs' handling of special characters in database and role names (Noah Misch, Nathan Bossart, Michael Paquier)

Numerous places in vacuumdb and other client programs could become confused by database and role names containing double quotes or backslashes. Tighten up quoting rules to make that safe. Also, ensure that when a conninfo string is used as a database name parameter to these programs, it is correctly treated as such throughout.

Fix handling of paired double quotes in psq|l's \connect and \password commands to match the documentation.

Introduce a new -reuse-previous option in psq|l's \connect command to allow explicit control of whether to re-use connection parameters from a previous connection. (Without this, the choice is based on whether the database name looks like a conninfo string, as before.) This allows secure handling of database names containing special characters in pg_dumpall scripts.

pg_dumpall now refuses to deal with database and role names containing carriage returns or new-lines, as it seems impractical to quote those characters safely on Windows. In future we may reject such names on the server side, but that step has not been taken yet.

These are considered security fixes because crafted object names containing special characters could have been used to execute commands with superuser privileges the next time a superuser executes pg_dumpall or other routine maintenance operations. (CVE-2016-5424)

• Fix corner-case misbehaviors for IS NULL/IS NOT NULL applied to nested composite values (Andrew Gierth, Tom Lane)

The SQL standard specifies that IS NULL should return TRUE for a row of all null values (thus ROW(NULL,NULL) IS NULL yields TRUE), but this is not meant to apply recursively (thus ROW(NULL, ROW(NULL,NULL)) IS NULL yields FALSE). The core executor got this right, but certain planner optimizations treated the test as recursive (thus producing TRUE in both cases), and contrib/postgres_fdw could produce remote queries that misbehaved similarly.

• Make the inet and cidr data types properly reject IPv6 addresses with too many colon-separated fields (Tom Lane)

• Prevent crash in close_ps() (the point #1 lseg operator) for NaN input coordinates (Tom Lane)

Make it return NULL instead of crashing.

• Fix several one-byte buffer over-reads in to_number() (Peter Eisentraut)
In several cases the `to_number()` function would read one more character than it should from the input string. There is a small chance of a crash, if the input happens to be adjacent to the end of memory.

• Avoid unsafe intermediate state during expensive paths through `heap_update()` (Masahiko Sawada, Andres Freund)

Previously, these cases locked the target tuple (by setting its XMAX) but did not WAL-log that action, thus risking data integrity problems if the page were spilled to disk and then a database crash occurred before the tuple update could be completed.

• Avoid crash in `postgres -C` when the specified variable has a null string value (Michael Paquier)

• Avoid consuming a transaction ID during `VACUUM` (Alexander Korotkov)

Some cases in `VACUUM` unnecessarily caused an XID to be assigned to the current transaction. Normally this is negligible, but if one is up against the XID wraparound limit, consuming more XIDs during anti-wraparound vacuums is a very bad thing.

• Avoid canceling hot-standby queries during `VACUUM FREEZE` (Simon Riggs, Álvaro Herrera)

`VACUUM FREEZE` on an otherwise-idle master server could result in unnecessary cancellations of queries on its standby servers.

• When a manual `ANALYZE` specifies a column list, don't reset the table's `changes_since_analyze` counter (Tom Lane)

If we're only analyzing some columns, we should not prevent routine auto-analyze from happening for the other columns.

• Fix `ANALYZE`'s overestimation of `n_distinct` for a unique or nearly-unique column with many null entries (Tom Lane)

The nulls could get counted as though they were themselves distinct values, leading to serious planner misestimates in some types of queries.

• Prevent autovacuum from starting multiple workers for the same shared catalog (Álvaro Herrera)

Normally this isn't much of a problem because the vacuum doesn't take long anyway; but in the case of a severely bloated catalog, it could result in all but one worker uselessly waiting instead of doing useful work on other tables.

• Prevent infinite loop in GiST index build for geometric columns containing NaN component values (Tom Lane)

• Fix `contrib/btree_gin` to handle the smallest possible `bigint` value correctly (Peter Eisentraut)

• Teach `libpq` to correctly decode server version from future servers (Peter Eisentraut)

It's planned to switch to two-part instead of three-part server version numbers for releases after 9.6. Make sure that `PQserverVersion()` returns the correct value for such cases.

• Fix ecpg's code for `unsigned long long` array elements (Michael Meskes)

• In `pg_dump` with both `-c` and `-C` options, avoid emitting an unwanted `CREATE SCHEMA public` command (David Johnston, Tom Lane)

• Make `pg_basebackup` accept `-Z 0` as specifying no compression (Fujii Masao)

• Fix `makefiles`' rule for building AIX shared libraries to be safe for parallel make (Noah Misch)
• Fix TAP tests and MSVC scripts to work when build directory's path name contains spaces (Michael Paquier, Kyotaro Horiguchi)

• Make regression tests safe for Danish and Welsh locales (Jeff Janes, Tom Lane)
  Change some test data that triggered the unusual sorting rules of these locales.

• Update our copy of the timezone code to match IANA’s tzcode release 2016c (Tom Lane)
  This is needed to cope with anticipated future changes in the time zone data files. It also fixes some corner-case bugs in coping with unusual time zones.

• Update time zone data files to tzdata release 2016f for DST law changes in Kemerovo and Novosibirsk, plus historical corrections for Azerbaijan, Belarus, and Morocco.

E.92. Release 9.2.17

Release date: 2016-05-12

This release contains a variety of fixes from 9.2.16. For information about new features in the 9.2 major release, see Section E.109.

E.92.1. Migration to Version 9.2.17

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.92.2. Changes

• Clear the OpenSSL error queue before OpenSSL calls, rather than assuming it’s clear already; and make sure we leave it clear afterwards (Peter Geoghegan, Dave Vitek, Peter Eisentraut)
  This change prevents problems when there are multiple connections using OpenSSL within a single process and not all the code involved follows the same rules for when to clear the error queue. Failures have been reported specifically when a client application uses SSL connections in libpq concurrently with SSL connections using the PHP, Python, or Ruby wrappers for OpenSSL. It’s possible for similar problems to arise within the server as well, if an extension module establishes an outgoing SSL connection.

• Fix “failed to build any n-way joins” planner error with a full join enclosed in the right-hand side of a left join (Tom Lane)

• Fix incorrect handling of equivalence-class tests in multilevel nestloop plans (Tom Lane)
  Given a three-or-more-way equivalence class of variables, such as \(X \cdot X = Y \cdot Y = Z \cdot Z\), it was possible for the planner to omit some of the tests needed to enforce that all the variables are actually equal, leading to join rows being output that didn’t satisfy the \(\text{WHERE}\) clauses. For various reasons, erroneous plans were seldom selected in practice, so that this bug has gone undetected for a long time.

• Fix possible misbehavior of TH, th, and Y, YYY format codes in \texttt{to_timestamp()} (Tom Lane)
  These could advance off the end of the input string, causing subsequent format codes to read garbage.

• Fix dumping of rules and views in which the \texttt{array} argument of a \texttt{value operator ANY (array)} construct is a sub-SELECT (Tom Lane)
• Make pg_regress use a startup timeout from the PGCTLTIMEDOUT environment variable, if that’s set (Tom Lane)

This is for consistency with a behavior recently added to pg_ctl; it eases automated testing on slow machines.

• Fix pg_upgrade to correctly restore extension membership for operator families containing only one operator class (Tom Lane)

In such a case, the operator family was restored into the new database, but it was no longer marked as part of the extension. This had no immediate ill effects, but would cause later pg_dump runs to emit output that would cause (harmless) errors on restore.

• Back-port 9.4-era memory-barrier code changes into 9.2 and 9.3 (Tom Lane)

These changes were not originally needed in pre-9.4 branches, but we recently back-patched a fix that expected the barrier code to work properly. Only IA64 (when using icc), HPPA, and Alpha platforms are affected.

• Reduce the number of SysV semaphores used by a build configured with --disable-spin-locks (Tom Lane)

• Rename internal function strtoi() to strtoint() to avoid conflict with a NetBSD library function (Thomas Munro)

• Fix reporting of errors from bind() and listen() system calls on Windows (Tom Lane)

• Reduce verbosity of compiler output when building with Microsoft Visual Studio (Christian Ullrich)

• Avoid possibly-unsafe use of Windows' FormatMessage() function (Christian Ullrich)

  Use the FORMAT_MESSAGE_IGNORE_INSERTS flag where appropriate. No live bug is known to exist here, but it seems like a good idea to be careful.

• Update time zone data files to tzdata release 2016d for DST law changes in Russia and Venezuela. There are new zone names Europe/Kirov and Asia/Tomsk to reflect the fact that these regions now have different time zone histories from adjacent regions.

E.93. Release 9.2.16

Release date: 2016-03-31

This release contains a variety of fixes from 9.2.15. For information about new features in the 9.2 major release, see Section E.109.

E.93.1. Migration to Version 9.2.16

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.93.2. Changes

• Fix incorrect handling of NULL index entries in indexed ROW() comparisons (Tom Lane)

An index search using a row comparison such as ROW(a, b) > ROW('x', 'y') would stop upon reaching a NULL entry in the b column, ignoring the fact that there might be non-NULL b values associated with later values of a.
• Avoid unlikely data-loss scenarios due to renaming files without adequate fsync() calls before and after (Michael Paquier, Tomas Vondra, Andres Freund)

• Correctly handle cases where pg_subtrans is close to XID wraparound during server startup (Jeff Janes)

• Fix corner-case crash due to trying to free localeconv() output strings more than once (Tom Lane)

• Fix parsing of affix files for ispell dictionaries (Tom Lane)

  The code could go wrong if the affix file contained any characters whose byte length changes during case-folding, for example I in Turkish UTF8 locales.

• Avoid use of sscanf() to parse ispell dictionary files (Artur Zakirov)

  This dodges a portability problem on FreeBSD-derived platforms (including macOS).

• Avoid a crash on old Windows versions (before 7SP1/2008R2SP1) with an AVX2-capable CPU and a Postgres build done with Visual Studio 2013 (Christian Ullrich)

  This is a workaround for a bug in Visual Studio 2013's runtime library, which Microsoft have stated they will not fix in that version.

• Fix psql's tab completion logic to handle multibyte characters properly (Kyotaro Horiguchi, Robert Haas)

• Fix psql's tab completion for SECURITY LABEL (Tom Lane)

  Pressing TAB after SECURITY LABEL might cause a crash or offering of inappropriate keywords.

• Make pg_ctl accept a wait timeout from the PGCTLTIMEOUT environment variable, if none is specified on the command line (Noah Misch)

  This eases testing of slower buildfarm members by allowing them to globally specify a longer-than-normal timeout for postmaster startup and shutdown.

• Fix incorrect test for Windows service status in pg_ctl (Manuel Mathar)

  The previous set of minor releases attempted to fix pg_ctl to properly determine whether to send log messages to Window's Event Log, but got the test backwards.

• Fix pgbench to correctly handle the combination of -C and -M prepared options (Tom Lane)

• In PL/Perl, properly translate empty Postgres arrays into empty Perl arrays (Alex Hunsaker)

• Make PL/Python cope with function names that aren't valid Python identifiers (Jim Nasby)

• Fix multiple mistakes in the statistics returned by contrib/pgstattuple's pgstatindex() function (Tom Lane)

• Remove dependency on psed in MSVC builds, since it's no longer provided by core Perl (Michael Paquier, Andrew Dunstan)

• Update time zone data files to tzdata release 2016c for DST law changes in Azerbaijan, Chile, Haiti, Palestine, and Russia (Altai, Astrakhan, Kirov, Sakhalin, Ulyanovsk regions), plus historical corrections for Lithuania, Moldova, and Russia (Kaliningrad, Samara, Volgograd).

E.94. Release 9.2.15

Release date: 2016-02-11
This release contains a variety of fixes from 9.2.14. For information about new features in the 9.2 major release, see Section E.109.

### E.94.1. Migration to Version 9.2.15

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

### E.94.2. Changes

- **Fix** infinite loops and buffer-overrun problems in regular expressions (Tom Lane)
  
  Very large character ranges in bracket expressions could cause infinite loops in some cases, and memory overwrites in other cases. (CVE-2016-0773)

- **Perform** an immediate shutdown if the `postmaster.pid` file is removed (Tom Lane)
  
  The postmaster now checks every minute or so that `postmaster.pid` is still there and still contains its own PID. If not, it performs an immediate shutdown, as though it had received SIGQUIT. The main motivation for this change is to ensure that failed buildfarm runs will get cleaned up without manual intervention; but it also serves to limit the bad effects if a DBA forcibly removes `postmaster.pid` and then starts a new postmaster.

- **In** `SERIALIZABLE` transaction isolation mode, serialization anomalies could be missed due to race conditions during insertions (Kevin Grittner, Thomas Munro)

- **Fix** failure to emit appropriate WAL records when doing `ALTER TABLE ... SET TABLESPACE` for unlogged relations (Michael Paquier, Andres Freund)
  
  Even though the relation's data is unlogged, the move must be logged or the relation will be inaccessible after a standby is promoted to master.

- **Fix** possible misinitialization of unlogged relations at the end of crash recovery (Andres Freund, Michael Paquier)

- **Fix** `ALTER COLUMN TYPE` to reconstruct inherited check constraints properly (Tom Lane)

- **Fix** `REASSIGN OWNED` to change ownership of composite types properly (Álvaro Herrera)

- **Fix** `REASSIGN OWNED` and `ALTER OWNER` to correctly update granted-permissions lists when changing owners of data types, foreign data wrappers, or foreign servers (Bruce Momjian, Álvaro Herrera)

- **Fix** `REASSIGN OWNED` to ignore foreign user mappings, rather than fail (Álvaro Herrera)

- **Add** more defenses against bad planner cost estimates for GIN index scans when the index's internal statistics are very out-of-date (Tom Lane)

- **Make** planner cope with hypothetical GIN indexes suggested by an index advisor plug-in (Julien Rouhaud)

- **Fix** dumping of whole-row Vars in `ROW()` and `VALUES()` lists (Tom Lane)

- **Fix** possible internal overflow in `numeric` division (Dean Rasheed)

- **Fix** enforcement of restrictions inside parentheses within regular expression lookahead constraints (Tom Lane)
  
  Lookahead constraints aren't allowed to contain backrefs, and parentheses within them are always considered non-capturing, according to the manual. However, the code failed to handle these cases properly inside a parenthesized subexpression, and would give unexpected results.
• Conversion of regular expressions to indexscan bounds could produce incorrect bounds from reg-
  exps containing lookahead constraints (Tom Lane)

• Fix regular-expression compiler to handle loops of constraint arcs (Tom Lane)

The code added for CVE-2007-4772 was both incomplete, in that it didn't handle loops involving
more than one state, and incorrect, in that it could cause assertion failures (though there seem to be
no bad consequences of that in a non-assert build). Multi-state loops would cause the compiler to
run until the query was canceled or it reached the too-many-states error condition.

• Improve memory-usage accounting in regular-expression compiler (Tom Lane)

This causes the code to emit “regular expression is too complex” errors in some cases that previously
used unreasonable amounts of time and memory.

• Improve performance of regular-expression compiler (Tom Lane)

• Make %h and %r escapes in log_line_prefix work for messages emitted due to log_con-
  nections (Tom Lane)

Previously, %h/%r started to work just after a new session had emitted the “connection received”
log message; now they work for that message too.

• On Windows, ensure the shared-memory mapping handle gets closed in child processes that don't
  need it (Tom Lane, Amit Kapila)

This oversight resulted in failure to recover from crashes whenever logging_collector is
turned on.

• Fix possible failure to detect socket EOF in non-blocking mode on Windows (Tom Lane)

It's not entirely clear whether this problem can happen in pre-9.5 branches, but if it did, the symptom
would be that a walsender process would wait indefinitely rather than noticing a loss of connection.

• Avoid leaking a token handle during SSPI authentication (Christian Ullrich)

• In psql, ensure that libreadline's idea of the screen size is updated when the terminal window size
  changes (Merlin Moncure)

Previously, libreadline did not notice if the window was resized during query output, leading to
strange behavior during later input of multiline queries.

• Fix psql's \det command to interpret its pattern argument the same way as other \d commands
  with potentially schema-qualified patterns do (Reece Hart)

• Avoid possible crash in psql's \c command when previous connection was via Unix socket and
  command specifies a new hostname and same username (Tom Lane)

• In pg_ctl start -w, test child process status directly rather than relying on heuristics (Tom
  Lane, Michael Paquier)

Previously, pg_ctl relied on an assumption that the new postmaster would always create post-
master.pid within five seconds. But that can fail on heavily-loaded systems, causing pg_ctl to
report incorrectly that the postmaster failed to start.

Except on Windows, this change also means that a pg_ctl start -w done immediately after
another such command will now reliably fail, whereas previously it would report success if done
within two seconds of the first command.

• In pg_ctl start -w, don't attempt to use a wildcard listen address to connect to the postmaster
  (Kondo Yuta)
On Windows, `pg_ctl` would fail to detect postmaster startup if `listen_addresses` is set to `0.0.0.0` or `::`, because it would try to use that value verbatim as the address to connect to, which doesn't work. Instead assume that `127.0.0.1` or `::1`, respectively, is the right thing to use.

- In `pg_ctl` on Windows, check service status to decide where to send output, rather than checking if standard output is a terminal (Michael Paquier)

- In `pg_dump` and `pg_basebackup`, adopt the GNU convention for handling tar-archive members exceeding 8GB (Tom Lane)

  The POSIX standard for `tar` file format does not allow archive member files to exceed 8GB, but most modern implementations of tar support an extension that fixes that. Adopt this extension so that `pg_dump` with `-Ft` no longer fails on tables with more than 8GB of data, and so that `pg_basebackup` can handle files larger than 8GB. In addition, fix some portability issues that could cause failures for members between 4GB and 8GB on some platforms. Potentially these problems could cause unrecoverable data loss due to unreadable backup files.

- Fix assorted corner-case bugs in `pg_dump`’s processing of extension member objects (Tom Lane)

- Make `pg_dump` mark a view’s triggers as needing to be processed after its rule, to prevent possible failure during parallel `pg_restore` (Tom Lane)

- Ensure that relation option values are properly quoted in `pg_dump` (Kouhei Sutou, Tom Lane)

  A reloption value that isn’t a simple identifier or number could lead to dump/reload failures due to syntax errors in `CREATE` statements issued by `pg_dump`. This is not an issue with any reloption currently supported by core PostgreSQL, but extensions could allow reloptions that cause the problem.

- Fix `pg_upgrade`’s file-copying code to handle errors properly on Windows (Bruce Momjian)

- Install guards in `pgbench` against corner-case overflow conditions during evaluation of script-specified division or modulo operators (Fabien Coelho, Michael Paquier)

- Fix failure to localize messages emitted by `pg_receivexlog` and `pg_recvlogical` (Ioseph Kim)

- Avoid dump/reload problems when using both `plpython2` and `plpython3` (Tom Lane)

  In principle, both versions of PL/Python can be used in the same database, though not in the same session (because the two versions of libpython cannot safely be used concurrently). However, `pg_restore` and `pg_upgrade` both do things that can fall foul of the same-session restriction. Work around that by changing the timing of the check.

- Fix PL/Python regression tests to pass with Python 3.5 (Peter Eisentraut)

- Prevent certain PL/Java parameters from being set by non-superusers (Noah Misch)

  This change mitigates a PL/Java security bug (CVE-2016-0766), which was fixed in PL/Java by marking these parameters as superuser-only. To fix the security hazard for sites that update PostgreSQL more frequently than PL/Java, make the core code aware of them also.

- Improve `libpq`’s handling of out-of-memory situations (Michael Paquier, Amit Kapila, Heikki Linnakangas)

- Fix order of arguments in `ecpg`-generated `typedef` statements (Michael Meskes)

- Use `%g` not `%f` format in `ecpg`’s `PGTYPESnumeric_from_double()` (Tom Lane)

- Fix `ecpg`-supplied header files to not contain comments continued from a preprocessor directive line onto the next line (Michael Meskes)
Such a comment is rejected by ecpg. It's not yet clear whether ecpg itself should be changed.

- Ensure that contrib/pgcrypto's crypt() function can be interrupted by query cancel (Andreas Karlsson)

- Accept flex versions later than 2.5.x (Tom Lane, Michael Paquier)

  Now that flex 2.6.0 has been released, the version checks in our build scripts needed to be adjusted.

- Install our missing script where PGXS builds can find it (Jim Nasby)

  This allows sane behavior in a PGXS build done on a machine where build tools such as bison are missing.

- Ensure that dynloader.h is included in the installed header files in MSVC builds (Bruce Momjian, Michael Paquier)

- Add variant regression test expected-output file to match behavior of current libxml2 (Tom Lane)

  The fix for libxml2's CVE-2015-7499 causes it not to output error context reports in some cases where it used to do so. This seems to be a bug, but we'll probably have to live with it for some time, so work around it.

- Update time zone data files to tzdata release 2016a for DST law changes in Cayman Islands, Metlakatla, and Trans-Baikal Territory (Zabaykalsky Krai), plus historical corrections for Pakistan.

E.95. Release 9.2.14

**Release date:** 2015-10-08

This release contains a variety of fixes from 9.2.13. For information about new features in the 9.2 major release, see Section E.109.

E.95.1. Migration to Version 9.2.14

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.95.2. Changes

- Fix contrib/pgcrypto to detect and report too-short crypt() salts (Josh Kupershmidt)

  Certain invalid salt arguments crashed the server or disclosed a few bytes of server memory. We have not ruled out the viability of attacks that arrange for presence of confidential information in the disclosed bytes, but they seem unlikely. (CVE-2015-5288)

- Fix subtransaction cleanup after a portal (cursor) belonging to an outer subtransaction fails (Tom Lane, Michael Paquier)

  A function executed in an outer-subtransaction cursor could cause an assertion failure or crash by referencing a relation created within an inner subtransaction.

- Fix insertion of relations into the relation cache "init file" (Tom Lane)

  An oversight in a patch in the most recent minor releases caused pg_trigger_tgrelid_tgname_index to be omitted from the init file. Subsequent sessions detected this, then deemed the
init file to be broken and silently ignored it, resulting in a significant degradation in session startup
time. In addition to fixing the bug, install some guards so that any similar future mistake will be
more obvious.

- Avoid O(N^2) behavior when inserting many tuples into a SPI query result (Neil Conway)
- Improve LISTEN startup time when there are many unread notifications (Matt Newell)
- Back-patch 9.3-era addition of per-resource-owner lock caches (Jeff Janes)

This substantially improves performance when pg_dump tries to dump a large number of tables.

- Disable SSL renegotiation by default (Michael Paquier, Andres Freund)

While use of SSL renegotiation is a good idea in theory, we have seen too many bugs in practice,
both in the underlying OpenSSL library and in our usage of it. Renegotiation will be removed en-
tirely in 9.5 and later. In the older branches, just change the default value of `ssl_renegotiation_limit` to zero (disabled).

- Lower the minimum values of the *_freeze_max_age parameters (Andres Freund)

This is mainly to make tests of related behavior less time-consuming, but it may also be of value
for installations with limited disk space.

- Limit the maximum value of `wal_buffers` to 2GB to avoid server crashes (Josh Berkus)
- Fix rare internal overflow in multiplication of numeric values (Dean Rasheed)
- Guard against hard-to-reach stack overflows involving record types, range types, `json`, `jsonb`,
  `tsquery`, `ltxtquery` and `query_int` (Noah Misch)
- Fix handling of `DOW` and `DOY` in datetime input (Greg Stark)

These tokens aren't meant to be used in datetime values, but previously they resulted in opaque
internal error messages rather than “invalid input syntax”.

- Add more query-cancel checks to regular expression matching (Tom Lane)
- Add recursion depth protections to regular expression, SIMILAR TO, and LIKE matching (Tom
  Lane)

Suitable search patterns and a low stack depth limit could lead to stack-overrun crashes.

- Fix potential infinite loop in regular expression execution (Tom Lane)

A search pattern that can apparently match a zero-length string, but actually doesn't match because
of a back reference, could lead to an infinite loop.

- In regular expression execution, correctly record match data for capturing parentheses within a
  quantifier even when the match is zero-length (Tom Lane)
- Fix low-memory failures in regular expression compilation (Andreas Seltenreich)
- Fix low-probability memory leak during regular expression execution (Tom Lane)
- Fix rare low-memory failure in lock cleanup during transaction abort (Tom Lane)
- Fix “unexpected out-of-memory situation during sort” errors when using tuplestores with small
  `work_mem` settings (Tom Lane)
- Fix very-low-probability stack overrun in `qsort` (Tom Lane)
• Fix “invalid memory alloc request size” failure in hash joins with large work_mem settings (Tomas Vondra, Tom Lane)

• Fix assorted planner bugs (Tom Lane)

These mistakes could lead to incorrect query plans that would give wrong answers, or to assertion failures in assert-enabled builds, or to odd planner errors such as “could not devise a query plan for the given query”, “could not find pathkey item to sort”, “plan should not reference subplan’s variable”, or “failed to assign all NestLoopParams to plan nodes”. Thanks are due to Andreas Seltenreich and Piotr Stefaniak for fuzz testing that exposed these problems.

• Improve planner’s performance for UPDATE/DELETE on large inheritance sets (Tom Lane, Dean Rasheed)

• Ensure standby promotion trigger files are removed at postmaster startup (Michael Paquier, Fujii Masao)

This prevents unwanted promotion from occurring if these files appear in a database backup that is used to initialize a new standby server.

• During postmaster shutdown, ensure that per-socket lock files are removed and listen sockets are closed before we remove the postmaster.pid file (Tom Lane)

This avoids race-condition failures if an external script attempts to start a new postmaster as soon as pg_ctl stop returns.

• Fix postmaster’s handling of a startup-process crash during crash recovery (Tom Lane)

If, during a crash recovery cycle, the startup process crashes without having restored database consistency, we’d try to launch a new startup process, which typically would just crash again, leading to an infinite loop.

• Do not print a WARNING when an autovacuum worker is already gone when we attempt to signal it, and reduce log verbosity for such signals (Tom Lane)

• Prevent autovacuum launcher from sleeping unduly long if the server clock is moved backwards a large amount (Álvaro Herrera)

• Ensure that cleanup of a GIN index's pending-insertions list is interruptable by cancel requests (Jeff Janes)

• Allow all-zeroes pages in GIN indexes to be reused (Heikki Linnakangas)

Such a page might be left behind after a crash.

• Fix handling of all-zeroes pages in SP-GiST indexes (Heikki Linnakangas)

VACUUM attempted to recycle such pages, but did so in a way that wasn't crash-safe.

• Fix off-by-one error that led to otherwise-harmless warnings about “apparent wraparound” in subtrans/multixact truncation (Thomas Munro)

• Fix misreporting of CONTINUE and MOVE statement types in PL/pgSQL’s error context messages (Pavel Stehule, Tom Lane)

• Fix PL/Perl to handle non-ASCII error message texts correctly (Alex Hunsaker)

• Fix PL/Python crash when returning the string representation of a record result (Tom Lane)

• Fix some places in PL/Tcl that neglected to check for failure of malloc() calls (Michael Paquier, Álvaro Herrera)
- In contrib/isn, fix output of ISBN-13 numbers that begin with 979 (Fabien Coelho)

  EANs beginning with 979 (but not 9790) are considered ISBNs, but they must be printed in the new
  13-digit format, not the 10-digit format.

- Fix contrib/sepgsql's handling of SELECT INTO statements (Kohei KaiGai)

- Improve libpq's handling of out-of-memory conditions (Michael Paquier, Heikki Linnakangas)

- Fix memory leaks and missing out-of-memory checks in ecpg (Michael Paquier)

- Fix psql's code for locale-aware formatting of numeric output (Tom Lane)

  The formatting code invoked by \pset numericlocale on did the wrong thing for some
  uncommon cases such as numbers with an exponent but no decimal point. It could also mangle
  already-localized output from the money data type.

- Prevent crash in psql's \c command when there is no current connection (Noah Misch)

- Make pg_dump handle inherited NOT VALID check constraints correctly (Tom Lane)

- Fix selection of default zlib compression level in pg_dump's directory output format (Andrew Dun-
  stan)

- Ensure that temporary files created during a pg_dump run with tar-format output are not world-
  readable (Michael Paquier)

- Fix pg_dump and pg_upgrade to support cases where the postgres or template1 database is
  in a non-default tablesapce (Marti Raudsepp, Bruce Momjian)

- Fix pg_dump to handle object privileges sanely when dumping from a server too old to have a
  particular privilege type (Tom Lane)

  When dumping data types from pre-9.2 servers, and when dumping functions or procedural lan-
  guages from pre-7.3 servers, pg_dump would produce GRANT/REVOKE commands that revoked
  the owner's grantable privileges and instead granted all privileges to PUBLIC. Since the privileges
  involved are just USAGE and EXECUTE, this isn't a security problem, but it's certainly a surprising
  representation of the older systems' behavior. Fix it to leave the default privilege state alone in these
  cases.

- Fix pg_dump to dump shell types (Tom Lane)

  Shell types (that is, not-yet-fully-defined types) aren't useful for much, but nonetheless pg_dump
  should dump them.

- Fix assorted minor memory leaks in pg_dump and other client-side programs (Michael Paquier)

- Fix spinlock assembly code for PPC hardware to be compatible with AIX's native assembler (Tom
  Lane)

  Building with gcc didn't work if gcc had been configured to use the native assembler, which is
  becoming more common.

- On AIX, test the -qlonglong compiler option rather than just assuming it's safe to use (Noah
  Misch)

- On AIX, use -Wl,-brtllib link option to allow symbols to be resolved at runtime (Noah Misch)

  Perl relies on this ability in 5.8.0 and later.

- Avoid use of inline functions when compiling with 32-bit xlc, due to compiler bugs (Noah Misch)
• Use `librt` for `sched_yield()` when necessary, which it is on some Solaris versions (Oskari Saarenmaa)

• Fix Windows `install.bat` script to handle target directory names that contain spaces (Heikki Linnakangas)

• Make the numeric form of the PostgreSQL version number (e.g., `90405`) readily available to extension Makefiles, as a variable named `VERSION_NUM` (Michael Paquier)

• Update time zone data files to `tzdata` release 2015g for DST law changes in Cayman Islands, Fiji, Moldova, Morocco, Norfolk Island, North Korea, Turkey, and Uruguay. There is a new zone name `America/Fort_Nelson` for the Canadian Northern Rockies.

E.96. Release 9.2.13

Release date: 2015-06-12

This release contains a small number of fixes from 9.2.12. For information about new features in the 9.2 major release, see Section E.109.

E.96.1. Migration to Version 9.2.13

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.96.2. Changes

• Fix rare failure to invalidate relation cache init file (Tom Lane)

  With just the wrong timing of concurrent activity, a `VACUUM FULL` on a system catalog might fail to update the “init file” that's used to avoid cache-loading work for new sessions. This would result in later sessions being unable to access that catalog at all. This is a very ancient bug, but it's so hard to trigger that no reproducible case had been seen until recently.

• Avoid deadlock between incoming sessions and `CREATE/DROP DATABASE` (Tom Lane)

  A new session starting in a database that is the target of a `DROP DATABASE` command, or is the template for a `CREATE DATABASE` command, could cause the command to wait for five seconds and then fail, even if the new session would have exited before that.


Release date: 2015-06-04

This release contains a small number of fixes from 9.2.11. For information about new features in the 9.2 major release, see Section E.109.

E.97.1. Migration to Version 9.2.12

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.11, see Section E.98.

E.97.2. Changes

• Avoid failures while `fsync`ing data directory during crash restart (Abhijit Menon-Sen, Tom Lane)
In the previous minor releases we added a patch to _fsync_ everything in the data directory after a crash. Unfortunately its response to any error condition was to fail, thereby preventing the server from starting up, even when the problem was quite harmless. An example is that an unwritable file in the data directory would prevent restart on some platforms, but it is common to make SSL certificate files unwritable by the server. Revise this behavior so that permissions failures are ignored altogether, and other types of failures are logged but do not prevent continuing.

- Fix `pg_get_functiondef()` to show functions' LEAKPROOF property, if set (Jeevan Chalke)
- Remove configure's check prohibiting linking to a threaded libpython on OpenBSD (Tom Lane)

The failure this restriction was meant to prevent seems to not be a problem anymore on current OpenBSD versions.

- Allow libpq to use TLS protocol versions beyond v1 (Noah Misch)

For a long time, libpq was coded so that the only SSL protocol it would allow was TLS v1. Now that newer TLS versions are becoming popular, allow it to negotiate the highest commonly-supported TLS version with the server. (PostgreSQL servers were already capable of such negotiation, so no change is needed on the server side.) This is a back-patch of a change already released in 9.4.0.

E.98. Release 9.2.11

**Release date:** 2015-05-22

This release contains a variety of fixes from 9.2.10. For information about new features in the 9.2 major release, see Section E.109.

E.98.1. Migration to Version 9.2.11

A dump/restore is not required for those running 9.2.X.

However, if you use `contrib/citext`'s `regexp_matches()` functions, see the changelog entry below about that.

Also, if you are upgrading from a version earlier than 9.2.10, see Section E.99.

E.98.2. Changes

- Avoid possible crash when client disconnects just before the authentication timeout expires (Benkocs Norbert Attila)

  If the timeout interrupt fired partway through the session shutdown sequence, SSL-related state would be freed twice, typically causing a crash and hence denial of service to other sessions. Experimentation shows that an unauthenticated remote attacker could trigger the bug somewhat consistently, hence treat as security issue. (CVE-2015-3165)

- Improve detection of system-call failures (Noah Misch)

  Our replacement implementation of `snprintf()` failed to check for errors reported by the underlying system library calls; the main case that might be missed is out-of-memory situations. In the worst case this might lead to information exposure, due to our code assuming that a buffer had been overwritten when it hadn't been. Also, there were a few places in which security-relevant calls of other system library functions did not check for failure.

  It remains possible that some calls of the `printf()` family of functions are vulnerable to information disclosure if an out-of-memory error occurs at just the wrong time. We judge the risk to not be large, but will continue analysis in this area. (CVE-2015-3166)
• In contrib/pgcrypto, uniformly report decryption failures as “Wrong key or corrupt data” (Noah Misch)

Previously, some cases of decryption with an incorrect key could report other error message texts. It has been shown that such variance in error reports can aid attackers in recovering keys from other systems. While it’s unknown whether pgcrypto’s specific behaviors are likewise exploitable, it seems better to avoid the risk by using a one-size-fits-all message. (CVE-2015-3167)

• Fix incorrect declaration of contrib/citext's regexp_matches() functions (Tom Lane)

These functions should return setof text[], like the core functions they are wrappers for; but they were incorrectly declared as returning just text[]. This mistake had two results: first, if there was no match you got a scalar null result, whereas what you should get is an empty set (zero rows). Second, the g flag was effectively ignored, since you would get only one result array even if there were multiple matches.

While the latter behavior is clearly a bug, there might be applications depending on the former behavior; therefore the function declarations will not be changed by default until PostgreSQL 9.5. In pre-9.5 branches, the old behavior exists in version 1.0 of the citext extension, while we have provided corrected declarations in version 1.1 (which is not installed by default). To adopt the fix in pre-9.5 branches, execute ALTER EXTENSION citext UPDATE TO '1.1' in each database in which citext is installed. (You can also “update” back to 1.0 if you need to undo that.) Be aware that either update direction will require dropping and recreating any views or rules that use citext's regexp_matches() functions.

• Fix incorrect checking of deferred exclusion constraints after a HOT update (Tom Lane)

If a new row that potentially violates a deferred exclusion constraint is HOT-updated (that is, no indexed columns change and the row can be stored back onto the same table page) later in the same transaction, the exclusion constraint would be reported as violated when the check finally occurred, even if the row(s) the new row originally conflicted with had been deleted.

• Fix planning of star-schema-style queries (Tom Lane)

Sometimes, efficient scanning of a large table requires that index parameters be provided from more than one other table (commonly, dimension tables whose keys are needed to index a large fact table). The planner should be able to find such plans, but an overly restrictive search heuristic prevented it.

• Prevent improper reordering of antijoins (NOT EXISTS joins) versus other outer joins (Tom Lane)

This oversight in the planner has been observed to cause “could not find RelOptInfo for given relids” errors, but it seems possible that sometimes an incorrect query plan might get past that consistency check and result in silently-wrong query output.

• Fix incorrect matching of subexpressions in outer-join plan nodes (Tom Lane)

Previously, if textually identical non-strict subexpressions were used both above and below an outer join, the planner might try to re-use the value computed below the join, which would be incorrect because the executor would force the value to NULL in case of an unmatched outer row.

• Fix GEQO planner to cope with failure of its join order heuristic (Tom Lane)

This oversight has been seen to lead to “failed to join all relations together” errors in queries involving LATERAL, and that might happen in other cases as well.

• Fix possible deadlock at startup when max_prepared_transactions is too small (Heikki Linnakangas)

• Don’t archive useless preallocated WAL files after a timeline switch (Heikki Linnakangas)

• Avoid “cannot GetMultiXactIdMembers() during recovery” error (Álvaro Herrera)
• Recursively \texttt{fsync()} the data directory after a crash (Abhijit Menon-Sen, Robert Haas)

This ensures consistency if another crash occurs shortly later. (The second crash would have to be a system-level crash, not just a database crash, for there to be a problem.)

• Fix autovacuum launcher’s possible failure to shut down, if an error occurs after it receives SIGTERM (Álvaro Herrera)

• Cope with unexpected signals in \texttt{LockBufferForCleanup()} (Andres Freund)

This oversight could result in spurious errors about “multiple backends attempting to wait for pin-count 1”.

• Fix crash when doing \texttt{COPY IN} to a table with check constraints that contain whole-row references (Tom Lane)

The known failure case only crashes in 9.4 and up, but there is very similar code in 9.3 and 9.2, so back-patch those branches as well.

• Avoid waiting for WAL flush or synchronous replication during commit of a transaction that was read-only so far as the user is concerned (Andres Freund)

Previously, a delay could occur at commit in transactions that had written WAL due to HOT page pruning, leading to undesirable effects such as sessions getting stuck at startup if all synchronous replicas are down. Sessions have also been observed to get stuck in catchup interrupt processing when using synchronous replication; this will fix that problem as well.

• Fix crash when manipulating hash indexes on temporary tables (Heikki Linnakangas)

• Fix possible failure during hash index bucket split, if other processes are modifying the index concurrently (Tom Lane)

• Check for interrupts while analyzing index expressions (Jeff Janes)

\texttt{ANALYZE} executes index expressions many times; if there are slow functions in such an expression, it’s desirable to be able to cancel the \texttt{ANALYZE} before that loop finishes.

• Ensure \texttt{tableoid} of a foreign table is reported correctly when a \texttt{READ COMMITTED} recheck occurs after locking rows in \texttt{SELECT FOR UPDATE, UPDATE, or DELETE} (Etsuro Fujita)

• Add the name of the target server to object description strings for foreign-server user mappings (Álvaro Herrera)

• Recommend setting \texttt{include_realm} to 1 when using Kerberos/GSSAPI/SSPI authentication (Stephen Frost)

Without this, identically-named users from different realms cannot be distinguished. For the moment this is only a documentation change, but it will become the default setting in PostgreSQL 9.5.

• Remove code for matching IPv4 \texttt{pg_hba.conf} entries to IPv4-in-IPv6 addresses (Tom Lane)

This hack was added in 2003 in response to a report that some Linux kernels of the time would report IPv4 connections as having IPv4-in-IPv6 addresses. However, the logic was accidentally broken in 9.0. The lack of any field complaints since then shows that it’s not needed anymore. Now we have reports that the broken code causes crashes on some systems, so let’s just remove it rather than fix it. (Had we chosen to fix it, that would make for a subtle and potentially security-sensitive change in the effective meaning of IPv4 \texttt{pg_hba.conf} entries, which does not seem like a good thing to do in minor releases.)

• Report WAL flush, not insert, position in \texttt{IDENTIFY_SYSTEM} replication command (Heikki Linnakangas)
This avoids a possible startup failure in pg_receivexlog.

- While shutting down service on Windows, periodically send status updates to the Service Control Manager to prevent it from killing the service too soon; and ensure that pg_ctl will wait for shutdown (Krystian Bigaj)

- Reduce risk of network deadlock when using libpq's non-blocking mode (Heikki Linnakangas)

  When sending large volumes of data, it's important to drain the input buffer every so often, in case the server has sent enough response data to cause it to block on output. (A typical scenario is that the server is sending a stream of NOTICE messages during \texttt{COPY FROM STDIN}.) This worked properly in the normal blocking mode, but not so much in non-blocking mode. We've modified libpq to opportunistically drain input when it can, but a full defense against this problem requires application cooperation: the application should watch for socket read-ready as well as write-ready conditions, and be sure to call \texttt{PQconsumeInput()} upon read-ready.

- In libpq, fix misparsing of empty values in URI connection strings (Thomas Fanghaenel)

- Fix array handling in ecpg (Michael Meskes)

- Fix psql to sanely handle URIs and conninfo strings as the first parameter to \texttt{\textbackslash connect} (David Fetter, Andrew Dunstan, Álvaro Herrera)

  This syntax has been accepted (but undocumented) for a long time, but previously some parameters might be taken from the old connection instead of the given string, which was agreed to be undesirable.

- Suppress incorrect complaints from psql on some platforms that it failed to write \texttt{~/.psql_history} at exit (Tom Lane)

  This misbehavior was caused by a workaround for a bug in very old (pre-2006) versions of libedit. We fixed it by removing the workaround, which will cause a similar failure to appear for anyone still using such versions of libedit. Recommendation: upgrade that library, or use libreadline.

- Fix \texttt{pg_dump}'s rule for deciding which casts are system-provided casts that should not be dumped (Tom Lane)

- In \texttt{pg_dump}, fix failure to honor \texttt{-Z} compression level option together with \texttt{-Fd} (Michael Paquier)

- Make \texttt{pg_dump} consider foreign key relationships between extension configuration tables while choosing dump order (Gilles Darold, Michael Paquier, Stephen Frost)

  This oversight could result in producing dumps that fail to reload because foreign key constraints are transiently violated.

- Fix dumping of views that are just \texttt{VALUES(...)} but have column aliases (Tom Lane)

- In \texttt{pg_upgrade}, force timeline 1 in the new cluster (Bruce Momjian)

  This change prevents upgrade failures caused by bogus complaints about missing WAL history files.

- In \texttt{pg_upgrade}, check for improperly non-connectable databases before proceeding (Bruce Momjian)

- In \texttt{pg_upgrade}, quote directory paths properly in the generated \texttt{delete_old_cluster} script (Bruce Momjian)

- In \texttt{pg_upgrade}, preserve database-level freezing info properly (Bruce Momjian)

  This oversight could cause missing-clog-file errors for tables within the \texttt{postgres} and template1 databases.
• Run pg_upgrade and pg_resetxlog with restricted privileges on Windows, so that they don't fail when run by an administrator (Muhammad Asif Naeem)

• Improve handling of readdir() failures when scanning directories in initdb and pg_basebackup (Marco Nenciarini)

• Fix failure in pg_receivexlog (Andres Freund)

  A patch merge mistake in 9.2.10 led to “could not create archive status file” errors.

• Fix slow sorting algorithm in contrib/intarray (Tom Lane)

• Fix compile failure on Sparc V8 machines (Rob Rowan)

• Update time zone data files to tzdata release 2015d for DST law changes in Egypt, Mongolia, and Palestine, plus historical changes in Canada and Chile. Also adopt revised zone abbreviations for the America/Adak zone (HST/HDT not HAST/HADT).

E.99. Release 9.2.10

Release date: 2015-02-05

This release contains a variety of fixes from 9.2.9. For information about new features in the 9.2 major release, see Section E.109.

E.99.1. Migration to Version 9.2.10

A dump/restore is not required for those running 9.2.X.

However, if you are a Windows user and are using the “Norwegian (Bokmål)” locale, manual action is needed after the upgrade to replace any “Norwegian (Bokmål)_Norway” locale names stored in PostgreSQL system catalogs with the plain-ASCII alias “Norwegian_Norway”. For details see http://wiki.postgresql.org/wiki/Changes_To_Norwegian_Locale

Also, if you are upgrading from a version earlier than 9.2.9, see Section E.100.

E.99.2. Changes

• Fix buffer overruns in to_char() (Bruce Momjian)

  When to_char() processes a numeric formatting template calling for a large number of digits, PostgreSQL would read past the end of a buffer. When processing a crafted timestamp formatting template, PostgreSQL would write past the end of a buffer. Either case could crash the server. We have not ruled out the possibility of attacks that lead to privilege escalation, though they seem unlikely. (CVE-2015-0241)

• Fix buffer overrun in replacement *printf() functions (Tom Lane)

  PostgreSQL includes a replacement implementation of printf and related functions. This code will overrun a stack buffer when formatting a floating point number (conversion specifiers e, E, f, F, g or G) with requested precision greater than about 500. This will crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. A database user can trigger such a buffer overrun through the to_char() SQL function. While that is the only affected core PostgreSQL functionality, extension modules that use printf-family functions may be at risk as well.

  This issue primarily affects PostgreSQL on Windows. PostgreSQL uses the system implementation of these functions where adequate, which it is on other modern platforms. (CVE-2015-0242)

• Fix buffer overruns in contrib/pgcrypto (Marko Tiikkaja, Noah Misch)
Errors in memory size tracking within the `pgcrypto` module permitted stack buffer overruns and improper dependence on the contents of uninitialized memory. The buffer overrun cases can crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. (CVE-2015-0243)

- **Fix possible loss of frontend/backend protocol synchronization after an error (Heikki Linnakangas)**

  If any error occurred while the server was in the middle of reading a protocol message from the client, it could lose synchronization and incorrectly try to interpret part of the message's data as a new protocol message. An attacker able to submit crafted binary data within a command parameter might succeed in injecting his own SQL commands this way. Statement timeout and query cancellation are the most likely sources of errors triggering this scenario. Particularly vulnerable are applications that use a timeout and also submit arbitrary user-crafted data as binary query parameters. Disabling statement timeout will reduce, but not eliminate, the risk of exploit. Our thanks to Emil Lenngren for reporting this issue. (CVE-2015-0244)

- **Fix information leak via constraint-violation error messages (Stephen Frost)**

  Some server error messages show the values of columns that violate a constraint, such as a unique constraint. If the user does not have SELECT privilege on all columns of the table, this could mean exposing values that the user should not be able to see. Adjust the code so that values are displayed only when they came from the SQL command or could be selected by the user. (CVE-2014-8161)

- **Lock down regression testing's temporary installations on Windows (Noah Misch)**

  Use SSPI authentication to allow connections only from the OS user who launched the test suite. This closes on Windows the same vulnerability previously closed on other platforms, namely that other users might be able to connect to the test postmaster. (CVE-2014-0067)

- **Cope with the Windows locale named “Norwegian (Bokmål)” (Heikki Linnakangas)**

  Non-ASCII locale names are problematic since it's not clear what encoding they should be represented in. Map the troublesome locale name to a plain-ASCII alias, “Norwegian_Norway”.

- **Avoid possible data corruption if `ALTER DATABASE SET TABLESPACE` is used to move a database to a new tablespace and then shortly later move it back to its original tablespace (Tom Lane)**

- **Avoid corrupting tables when `ANALYZE` inside a transaction is rolled back (Andres Freund, Tom Lane, Michael Paquier)**

  If the failing transaction had earlier removed the last index, rule, or trigger from the table, the table would be left in a corrupted state with the relevant `pg_class` flags not set though they should be.

- **Ensure that unlogged tables are copied correctly during `CREATE DATABASE` or `ALTER DATABASE SET TABLESPACE` (Pavan Deolasee, Andres Freund)**

- **Fix `DROP`’s dependency searching to correctly handle the case where a table column is recursively visited before its table (Petr Jelinek, Tom Lane)**

  This case is only known to arise when an extension creates both a datatype and a table using that datatype. The faulty code might refuse a `DROP EXTENSION` unless `CASCADE` is specified, which should not be required.

- **Fix use-of-already-freed-memory problem in EvalPlanQual processing (Tom Lane)**

  In `READ COMMITTED` mode, queries that lock or update recently-updated rows could crash as a result of this bug.

- **Fix planning of `SELECT FOR UPDATE` when using a partial index on a child table (Kyotaro Horiguchi)**
In **READ COMMITTED** mode, **SELECT FOR UPDATE** must also recheck the partial index's **WHERE** condition when rechecking a recently-updated row to see if it still satisfies the query's **WHERE** condition. This requirement was missed if the index belonged to an inheritance child table, so that it was possible to incorrectly return rows that no longer satisfy the query condition.

- Fix corner case wherein **SELECT FOR UPDATE** could return a row twice, and possibly miss returning other rows (Tom Lane)

In **READ COMMITTED** mode, a **SELECT FOR UPDATE** that is scanning an inheritance tree could incorrectly return a row from a prior child table instead of the one it should return from a later child table.

- Reject duplicate column names in the referenced-columns list of a **FOREIGN KEY** declaration (David Rowley)

This restriction is per SQL standard. Previously we did not reject the case explicitly, but later on the code would fail with bizarre-looking errors.

- Restore previous behavior of conversion of domains to JSON (Tom Lane)

This change causes domains over numeric and boolean to be treated like their base types for purposes of conversion to JSON. It worked like that before 9.3.5 and 9.2.9, but was unintentionally changed while fixing a related problem.

- Fix bugs in raising a **numeric** value to a large integral power (Tom Lane)

The previous code could get a wrong answer, or consume excessive amounts of time and memory before realizing that the answer must overflow.

- In **numeric_recv()**, truncate away any fractional digits that would be hidden according to the value's dscale field (Tom Lane)

A **numeric** value's display scale (dscale) should never be less than the number of nonzero fractional digits; but apparently there's at least one broken client application that transmits binary **numeric** values in which that's true. This leads to strange behavior since the extra digits are taken into account by arithmetic operations even though they aren't printed. The least risky fix seems to be to truncate away such “hidden” digits on receipt, so that the value is indeed what it prints as.

- Fix incorrect search for shortest-first regular expression matches (Tom Lane)

Matching would often fail when the number of allowed iterations is limited by a ? quantifier or a bound expression.

- Reject out-of-range numeric timezone specifications (Tom Lane)

Simple numeric timezone specifications exceeding +/- 168 hours (one week) would be accepted, but could then cause null-pointer dereference crashes in certain operations. There's no use-case for such large UTC offsets, so reject them.

- Fix bugs in tsquery @> tsquery operator (Heikki Linnakangas)

Two different terms would be considered to match if they had the same CRC. Also, if the second operand had more terms than the first, it would be assumed not to be contained in the first; which is wrong since it might contain duplicate terms.

- Improve ispell dictionary's defenses against bad affix files (Tom Lane)

- Allow more than 64K phrases in a thesaurus dictionary (David Boutin)

The previous coding could crash on an oversize dictionary, so this was deemed a back-patchable bug fix rather than a feature addition.
• Fix namespace handling in xpath() (Ali Akbar)

Previously, the xml value resulting from an xpath() call would not have namespace declarations if the namespace declarations were attached to an ancestor element in the input xml value, rather than to the specific element being returned. Propagate the ancestral declaration so that the result is correct when considered in isolation.

• Ensure that whole-row variables expose nonempty column names to functions that pay attention to column names within composite arguments (Tom Lane)

In some contexts, constructs like row_to_json(tab.*) may not produce the expected column names. This is fixed properly as of 9.4; in older branches, just ensure that we produce some non-empty name. (In some cases this will be the underlying table's column name rather than the query-assigned alias that should theoretically be visible.)

• Fix mishandling of system columns, particularly tableoid, in FDW queries (Etsuro Fujita)

• Avoid doing indexed_column = ANY (array) as an index qualifier if that leads to an inferior plan (Andrew Gierth)

In some cases, = ANY conditions applied to non-first index columns would be done as index conditions even though it would be better to use them as simple filter conditions.

• Fix planner problems with nested append relations, such as inherited tables within UNION ALL subqueries (Tom Lane)

• Fail cleanly when a GiST index tuple doesn't fit on a page, rather than going into infinite recursion (Andrew Gierth)

• Exempt tables that have per-table cost_limit and/or cost_delay settings from autovacuum's global cost balancing rules (Álvaro Herrera)

The previous behavior resulted in basically ignoring these per-table settings, which was unintended. Now, a table having such settings will be vacuumed using those settings, independently of what is going on in other autovacuum workers. This may result in heavier total I/O load than before, so such settings should be re-examined for sanity.

• Avoid wholesale autovacuuming when autovacuum is nominally off (Tom Lane)

Even when autovacuum is nominally off, we will still launch autovacuum worker processes to vacuum tables that are at risk of XID wraparound. However, such a worker process then proceeded to vacuum all tables in the target database, if they met the usual thresholds for autovacuuming. This is at best pretty unexpected; at worst it delays response to the wraparound threat. Fix it so that if autovacuum is turned off, workers only do anti-wraparound vacuums and not any other work.

• During crash recovery, ensure that unlogged relations are rewritten as empty and are synced to disk before recovery is considered complete (Abhijit Menon-Sen, Andres Freund)

This prevents scenarios in which unlogged relations might contain garbage data following database crash recovery.

• Fix race condition between hot standby queries and replaying a full-page image (Heikki Linnakan-gas)

This mistake could result in transient errors in queries being executed in hot standby.

• Fix several cases where recovery logic improperly ignored WAL records for COMMIT/ABORT PREPARED (Heikki Linnakangas)

The most notable oversight was that recovery_target_xid could not be used to stop at a two-phase commit.
• Prevent latest WAL file from being archived a second time at completion of crash recovery (Fujii Masao)

• Avoid creating unnecessary .ready marker files for timeline history files (Fujii Masao)

• Fix possible null pointer dereference when an empty prepared statement is used and the log_statement setting is mod or ddl (Fujii Masao)

• Change “pgstat wait timeout” warning message to be LOG level, and rephrase it to be more understandable (Tom Lane)

  This message was originally thought to be essentially a can’t-happen case, but it occurs often enough on our slower buildfarm members to be a nuisance. Reduce it to LOG level, and expend a bit more effort on the wording: it now reads “using stale statistics instead of current ones because stats collector is not responding”.

• Fix SPARC spinlock implementation to ensure correctness if the CPU is being run in a non-TSO coherency mode, as some non-Solaris kernels do (Andres Freund)

• Warn if macOS’s setlocale() starts an unwanted extra thread inside the postmaster (Noah Misch)

• Fix processing of repeated dbname parameters in PQconnectdbParams() (Alex Shulgin)

  Unexpected behavior ensued if the first occurrence of dbname contained a connection string or URI to be expanded.

• Ensure that libpq reports a suitable error message on unexpected socket EOF (Marko Tiikkaja, Tom Lane)

  Depending on kernel behavior, libpq might return an empty error string rather than something useful when the server unexpectedly closed the socket.

• Clear any old error message during PQreset() (Heikki Linnakangas)

  If PQreset() is called repeatedly, and the connection cannot be re-established, error messages from the failed connection attempts kept accumulating in the PGconn’s error string.

• Properly handle out-of-memory conditions while parsing connection options in libpq (Alex Shulgin, Heikki Linnakangas)

• Fix array overrun in ecpg’s version of ParseDateTime() (Michael Paquier)

• In initdb, give a clearer error message if a password file is specified but is empty (Mats Erik Andersson)

• Fix psql’s \s command to work nicely with libedit, and add pager support (Stepan Rutz, Tom Lane)

  When using libedit rather than readline, \s printed the command history in a fairly unreadable encoded format, and on recent libedit versions might fail altogether. Fix that by printing the history ourselves rather than having the library do it. A pleasant side-effect is that the pager is used if appropriate.

  This patch also fixes a bug that caused newline encoding to be applied inconsistently when saving the command history with libedit. Multiline history entries written by older psql versions will be read cleanly with this patch, but perhaps not vice versa, depending on the exact libedit versions involved.

• Improve consistency of parsing of psql’s special variables (Tom Lane)

  Allow variant spellings of on and off (such as 1/0) for ECHO_HIDDEN and ON_ERROR_Rollback. Report a warning for unrecognized values for COMP_KEYWORD_CASE, ECHO,
Release Notes

ECHO_HIDDEN, HISTCONTROL, ON_ERROR_ROLLBACK, and VERBOSITY. Recognize all values for all these variables case-insensitively; previously there was a mishmash of case-sensitive and case-insensitive behaviors.

- Fix psql's expanded-mode display to work consistently when using border = 3 and linestyle = ascii or unicode (Stephen Frost)
- Improve performance of pg_dump when the database contains many instances of multiple dependency paths between the same two objects (Tom Lane)
- Fix pg_dumpall to restore its ability to dump from pre-8.1 servers (Gilles Darold)
- Fix possible deadlock during parallel restore of a schema-only dump (Robert Haas, Tom Lane)
- Fix core dump in pg_dump --binary-upgrade on zero-column composite type (Rushabh Lathia)
- Prevent WAL files created by pg_basebackup -x/-X from being archived again when the standby is promoted (Andres Freund)
- Fix failure of contrib/auto_explain to print per-node timing information when doing EXPLAIN ANALYZE (Tom Lane)
- Fix upgrade-from-unpackaged script for contrib/citext (Tom Lane)
- Fix block number checking in contrib/pageinspect's get_raw_page() (Tom Lane)
  The incorrect checking logic could prevent access to some pages in non-main relation forks.
- Fix contrib/pgcrypto's pgp_sym_decrypt() to not fail on messages whose length is 6 less than a power of 2 (Marko Tiikkaja)
- Fix file descriptor leak in contrib/pg_test_fsync (Jeff Janes)
  This could cause failure to remove temporary files on Windows.
- Handle unexpected query results, especially NULLs, safely in contrib/tablefunc's connectby() (Michael Paquier)
  connectby() previously crashed if it encountered a NULL key value. It now prints that row but doesn't recurse further.
- Avoid a possible crash in contrib/xml2's xslt_process() (Mark Simonetti)
  libxml2 seems to have an undocumented dependency on the order in which resources are freed; reorder our calls to avoid a crash.
- Mark some contrib I/O functions with correct volatility properties (Tom Lane)

  The previous over-conservative marking was immaterial in normal use, but could cause optimization problems or rejection of valid index expression definitions. Since the consequences are not large, we've just adjusted the function definitions in the extension modules' scripts, without changing version numbers.

- Numerous cleanups of warnings from Coverity static code analyzer (Andres Freund, Tatsuo Ishii, Marko Kreen, Tom Lane, Michael Paquier)

  These changes are mostly cosmetic but in some cases fix corner-case bugs, for example a crash rather than a proper error report after an out-of-memory failure. None are believed to represent security issues.
- Detect incompatible OpenLDAP versions during build (Noah Misch)

2731
With OpenLDAP versions 2.4.24 through 2.4.31, inclusive, PostgreSQL backends can crash at exit. Raise a warning during configure based on the compile-time OpenLDAP version number, and test the crashing scenario in the contrib/dblink regression test.

- In non-MSVC Windows builds, ensure libpq.dll is installed with execute permissions (Noah Misch)

- Make pg_regress remove any temporary installation it created upon successful exit (Tom Lane)

  This results in a very substantial reduction in disk space usage during make check-world, since that sequence involves creation of numerous temporary installations.

- Support time zone abbreviations that change UTC offset from time to time (Tom Lane)

  Previously, PostgreSQL assumed that the UTC offset associated with a time zone abbreviation (such as EST) never changes in the usage of any particular locale. However this assumption fails in the real world, so introduce the ability for a zone abbreviation to represent a UTC offset that sometimes changes. Update the zone abbreviation definition files to make use of this feature in timezone locales that have changed the UTC offset of their abbreviations since 1970 (according to the IANA timezone database). In such timezones, PostgreSQL will now associate the correct UTC offset with the abbreviation depending on the given date.

- Update time zone abbreviations lists (Tom Lane)

  Add CST (China Standard Time) to our lists. Remove references to ADT as “Arabia Daylight Time”, an abbreviation that's been out of use since 2007; therefore, claiming there is a conflict with “Atlantic Daylight Time” doesn't seem especially helpful. Fix entirely incorrect GMT offsets for CKT (Cook Islands), FJT, and FJST (Fiji); we didn't even have them on the proper side of the date line.

- Update time zone data files to tzdata release 2015a.

  The IANA timezone database has adopted abbreviations of the form A\w*ST/A\w*DT for all Australian time zones, reflecting what they believe to be current majority practice Down Under. These names do not conflict with usage elsewhere (other than ACST for Acre Summer Time, which has been in disuse since 1994). Accordingly, adopt these names into our “Default” timezone abbreviation set. The “Australia” abbreviation set now contains only CST, EAST, EST, SAST, SAT, and WST, all of which are thought to be mostly historical usage. Note that SAST has also been changed to be South Africa Standard Time in the “Default” abbreviation set.

  Also, add zone abbreviations SRET (Asia/Srednekolymsk) and XJT (Asia/Urumqi), and use WSST/WSDT for western Samoa. Also, there were DST law changes in Chile, Mexico, the Turks & Caicos Islands (America/Grand_Turk), and Fiji. There is a new zone Pacific/Bougainville for portions of Papua New Guinea. Also, numerous corrections for historical (pre-1970) time zone data.

**E.100. Release 9.2.9**

**Release date:** 2014-07-24

This release contains a variety of fixes from 9.2.8. For information about new features in the 9.2 major release, see Section E.109.

**E.100.1. Migration to Version 9.2.9**

A dump/restore is not required for those running 9.2.X.

However, this release corrects an index corruption problem in some GiST indexes. See the first changelog entry below to find out whether your installation has been affected and what steps you should take if so.
Also, if you are upgrading from a version earlier than 9.2.6, see Section E.103.

E.100.2. Changes

• Correctly initialize padding bytes in contrib/btree_gist indexes on bit columns (Heikki Linnakangas)
  
  This error could result in incorrect query results due to values that should compare equal not being seen as equal. Users with GiST indexes on bit or bit varying columns should REINDEX those indexes after installing this update.

• Protect against torn pages when deleting GIN list pages (Heikki Linnakangas)
  
  This fix prevents possible index corruption if a system crash occurs while the page update is being written to disk.

• Don't clear the right-link of a GiST index page while replaying updates from WAL (Heikki Linnakangas)
  
  This error could lead to transiently wrong answers from GiST index scans performed in Hot Standby.

• Fix corner-case infinite loop during insertion into an SP-GiST text index (Tom Lane)

• Fix feedback status when hot_standby_feedback is turned off on-the-fly (Simon Riggs)

• Fix possibly-incorrect cache invalidation during nested calls to ReceiveSharedInvalidMessages (Andres Freund)

• Fix planner's mishandling of nested PlaceHolderVars generated in nested-nestloop plans (Tom Lane)
  
  This oversight could result in "variable not found in subplan target lists" errors, or in silently wrong query results.

• Fix "could not find pathkey item to sort" planner failures with UNION ALL over subqueries reading from tables with inheritance children (Tom Lane)

• Don't assume a subquery's output is unique if there's a set-returning function in its targetlist (David Rowley)
  
  This oversight could lead to misoptimization of constructs like WHERE x IN (SELECT y, generate_series(1,10) FROM t GROUP BY y).

• Improve planner to drop constant-NULL inputs of AND/OR when possible (Tom Lane)
  
  This change fixes some cases where the more aggressive parameter substitution done by 9.2 and later can lead to a worse plan than older versions produced.

• Fix identification of input type category in to_json() and friends (Tom Lane)
  
  This is known to have led to inadequate quoting of money fields in the JSON result, and there may have been wrong results for other data types as well.

• Fix failure to detoast fields in composite elements of structured types (Tom Lane)
  
  This corrects cases where TOAST pointers could be copied into other tables without being dereferenced. If the original data is later deleted, it would lead to errors like "missing chunk number 0 for toast value ..." when the now-dangling pointer is used.

• Fix "record type has not been registered" failures with whole-row references to the output of Append plan nodes (Tom Lane)
• Fix possible crash when invoking a user-defined function while rewinding a cursor (Tom Lane)

• Fix query-lifespan memory leak while evaluating the arguments for a function in FROM (Tom Lane)

• Fix session-lifespan memory leaks in regular-expression processing (Tom Lane, Arthur O'Dwyer, Greg Stark)

• Fix data encoding error in hungarian.stop (Tom Lane)

• Prevent foreign tables from being created with OIDS when default_with_oids is true (Etsuro Fujita)

• Fix liveness checks for rows that were inserted in the current transaction and then deleted by a now-rolled-back subtransaction (Andres Freund)

This could cause problems (at least spurious warnings, and at worst an infinite loop) if CREATE INDEX or CLUSTER were done later in the same transaction.

• Clear pg_stat_activity.xact_start during PREPARE TRANSACTION (Andres Freund)

After the PREPARE, the originating session is no longer in a transaction, so it should not continue to display a transaction start time.

• Fix REASSIGN OWNED to not fail for text search objects (Álvaro Herrera)

• Block signals during postmaster startup (Tom Lane)

This ensures that the postmaster will properly clean up after itself if, for example, it receives SIGINT while still starting up.

• Fix client host name lookup when processing pg_hba.conf entries that specify host names instead of IP addresses (Tom Lane)

Ensure that reverse-DNS lookup failures are reported, instead of just silently not matching such entries. Also ensure that we make only one reverse-DNS lookup attempt per connection, not one per host name entry, which is what previously happened if the lookup attempts failed.

• Allow the root user to use postgres -C variable and postgres --describe-config (MauMau)

The prohibition on starting the server as root does not need to extend to these operations, and relaxing it prevents failure of pg_ctl in some scenarios.

• Secure Unix-domain sockets of temporary postmasters started during make check (Noah Misch)

Any local user able to access the socket file could connect as the server's bootstrap superuser, then proceed to execute arbitrary code as the operating-system user running the test, as we previously noted in CVE-2014-0067. This change defends against that risk by placing the server's socket in a temporary, mode 0700 subdirectory of /tmp. The hazard remains however on platforms where Unix sockets are not supported, notably Windows, because then the temporary postmaster must accept local TCP connections.

A useful side effect of this change is to simplify make check testing in builds that override DEFAULT_PGSOCKET_DIR. Popular non-default values like /var/run/postgresql are often not writable by the build user, requiring workarounds that will no longer be necessary.

• Fix tablespace creation WAL replay to work on Windows (MauMau)

• Fix detection of socket creation failures on Windows (Bruce Momjian)

• On Windows, allow new sessions to absorb values of PGC_BACKEND parameters (such as log_connections) from the configuration file (Amit Kapila)
Previously, if such a parameter were changed in the file post-startup, the change would have no effect.

• Properly quote executable path names on Windows (Nikhil Deshpande)

This oversight could cause initdb and pg_upgrade to fail on Windows, if the installation path contained both spaces and @ signs.

• Fix linking of libpython on macOS (Tom Lane)

The method we previously used can fail with the Python library supplied by Xcode 5.0 and later.

• Avoid buffer bloat in libpq when the server consistently sends data faster than the client can absorb it (Shin-ichi Morita, Tom Lane)

libpq could be coerced into enlarging its input buffer until it runs out of memory (which would be reported misleadingly as “lost synchronization with server”). Under ordinary circumstances it’s quite far-fetched that data could be continuously transmitted more quickly than the recv() loop can absorb it, but this has been observed when the client is artificially slowed by scheduler constraints.

• Ensure that LDAP lookup attempts in libpq time out as intended (Laurenz Albe)

• Fix ecpg to do the right thing when an array of char * is the target for a FETCH statement returning more than one row, as well as some other array-handling fixes (Ashutosh Bapat)

• Fix pg_restore's processing of old-style large object comments (Tom Lane)

A direct-to-database restore from an archive file generated by a pre-9.0 version of pg_dump would usually fail if the archive contained more than a few comments for large objects.

• Fix pg_upgrade for cases where the new server creates a TOAST table but the old version did not (Bruce Momjian)

This rare situation would manifest as “relation OID mismatch” errors.

• Prevent contrib/auto_explain from changing the output of a user's EXPLAIN (Tom Lane)

If auto_explain is active, it could cause an EXPLAIN (ANALYZE, TIMING OFF) command to nonetheless print timing information.

• Fix query-lifespan memory leak in contrib/dblink (MauMau, Joe Conway)

• In contrib/pgcrypto functions, ensure sensitive information is cleared from stack variables before returning (Marko Kreen)

• Prevent use of already-freed memory in contrib/pgstattuple's pgstat_heap() (Noah Misch)

• In contrib/uuid-ossp, cache the state of the OSSP UUID library across calls (Tom Lane)

This improves the efficiency of UUID generation and reduces the amount of entropy drawn from /dev/urandom, on platforms that have that.

• Update time zone data files to tzdata release 2014e for DST law changes in Crimea, Egypt, and Morocco.

E.101. Release 9.2.8

Release date: 2014-03-20
This release contains a variety of fixes from 9.2.7. For information about new features in the 9.2 major release, see Section E.109.

E.101.1. Migration to Version 9.2.8

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.6, see Section E.103.

E.101.2. Changes

• Restore GIN metapages unconditionally to avoid torn-page risk (Heikki Linnakangas)

Although this oversight could theoretically result in a corrupted index, it is unlikely to have caused any problems in practice, since the active part of a GIN metapage is smaller than a standard 512-byte disk sector.

• Avoid race condition in checking transaction commit status during receipt of a NOTIFY message (Marko Tiikkaja)

This prevents a scenario wherein a sufficiently fast client might respond to a notification before database updates made by the notifier have become visible to the recipient.

• Allow regular-expression operators to be terminated early by query cancel requests (Tom Lane)

This prevents scenarios wherein a pathological regular expression could lock up a server process uninterruptibly for a long time.

• Remove incorrect code that tried to allow OVERLAPS with single-element row arguments (Joshua Yanovski)

This code never worked correctly, and since the case is neither specified by the SQL standard nor documented, it seemed better to remove it than fix it.

• Avoid getting more than AccessShareLock when de-parsing a rule or view (Dean Rasheed)

This oversight resulted in pg_dump unexpectedly acquiring RowExclusiveLock locks on tables mentioned as the targets of INSERT/UPDATE/DELETE commands in rules. While usually harmless, that could interfere with concurrent transactions that tried to acquire, for example, ShareLock on those tables.

• Improve performance of index endpoint probes during planning (Tom Lane)

This change fixes a significant performance problem that occurred when there were many not-yet-committed rows at the end of the index, which is a common situation for indexes on sequentially-assigned values such as timestamps or sequence-generated identifiers.

• Fix walsender’s failure to shut down cleanly when client is pg_receivexlog (Fujii Masao)

• Check WAL level and hot standby parameters correctly when doing crash recovery that will be followed by archive recovery (Heikki Linnakangas)

• Fix test to see if hot standby connections can be allowed immediately after a crash (Heikki Linnakangas)

• Prevent interrupts while reporting non-ERROR messages (Tom Lane)

This guards against rare server-process freezeups due to recursive entry to syslog(), and perhaps other related problems.
• Fix memory leak in PL/Perl when returning a composite result, including multiple-OUT-parameter cases (Alex Hunsaker)

• Fix tracking of psql script line numbers during `\copy` from out-of-line data (Kumar Rajeev Rastogi, Amit Khandekar)

`\copy` ... from incremented the script file line number for each data line, even if the data was not coming from the script file. This mistake resulted in wrong line numbers being reported for any errors occurring later in the same script file.

• Prevent intermittent “could not reserve shared memory region” failures on recent Windows versions (MauMau)

• Update time zone data files to tzdata release 2014a for DST law changes in Fiji and Turkey, plus historical changes in Israel and Ukraine.

E.102. Release 9.2.7

Release date: 2014-02-20

This release contains a variety of fixes from 9.2.6. For information about new features in the 9.2 major release, see Section E.109.

E.102.1. Migration to Version 9.2.7

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.6, see Section E.103.

E.102.2. Changes

• Shore up `GRANT ... WITH ADMIN OPTION` restrictions (Noah Misch)

Granting a role without `ADMIN OPTION` is supposed to prevent the grantee from adding or removing members from the granted role, but this restriction was easily bypassed by doing `SET ROLE` first. The security impact is mostly that a role member can revoke the access of others, contrary to the wishes of his grantor. Unapproved role member additions are a lesser concern, since an uncooperative role member could provide most of his rights to others anyway by creating views or `SECURITY DEFINER` functions. (CVE-2014-0060)

• Prevent privilege escalation via manual calls to PL validator functions (Andres Freund)

The primary role of PL validator functions is to be called implicitly during `CREATE FUNCTION`, but they are also normal SQL functions that a user can call explicitly. Calling a validator on a function actually written in some other language was not checked for and could be exploited for privilege-escalation purposes. The fix involves adding a call to a privilege-checking function in each validator function. Non-core procedural languages will also need to make this change to their own validator functions, if any. (CVE-2014-0061)

• Avoid multiple name lookups during table and index DDL (Robert Haas, Andres Freund)

If the name lookups come to different conclusions due to concurrent activity, we might perform some parts of the DDL on a different table than other parts. At least in the case of `CREATE INDEX`, this can be used to cause the permissions checks to be performed against a different table than the index creation, allowing for a privilege escalation attack. (CVE-2014-0062)

• Prevent buffer overrun with long datetime strings (Noah Misch)

The `MAXDATELEN` constant was too small for the longest possible value of type `interval`, allowing a buffer overrun in `interval_out()`. Although the datetime input functions were more
careful about avoiding buffer overrun, the limit was short enough to cause them to reject some valid inputs, such as input containing a very long timezone name. The ecpg library contained these vulnerabilities along with some of its own. (CVE-2014-0063)

- Prevent buffer overrun due to integer overflow in size calculations (Noah Misch, Heikki Linnakangas)

Several functions, mostly type input functions, calculated an allocation size without checking for overflow. If overflow did occur, a too-small buffer would be allocated and then written past. (CVE-2014-0064)

- Prevent overruns of fixed-size buffers (Peter Eisentraut, Jozef Mlich)

Use `strlcpy()` and related functions to provide a clear guarantee that fixed-size buffers are not overrun. Unlike the preceding items, it is unclear whether these cases really represent live issues, since in most cases there appear to be previous constraints on the size of the input string. Nonetheless it seems prudent to silence all Coverity warnings of this type. (CVE-2014-0065)

- Avoid crashing if `crypt()` returns NULL (Honza Horak, Bruce Momjian)

There are relatively few scenarios in which `crypt()` could return NULL, but `contrib/chkpass` would crash if it did. One practical case in which this could be an issue is if libc is configured to refuse to execute unapproved hashing algorithms (e.g., “FIPS mode”). (CVE-2014-0066)

- Document risks of `make check` in the regression testing instructions (Noah Misch, Tom Lane)

Since the temporary server started by `make check` uses “trust” authentication, another user on the same machine could connect to it as database superuser, and then potentially exploit the privileges of the operating-system user who started the tests. A future release will probably incorporate changes in the testing procedure to prevent this risk, but some public discussion is needed first. So for the moment, just warn people against using `make check` when there are untrusted users on the same machine. (CVE-2014-0067)

- Fix possible mis-replay of WAL records when some segments of a relation aren't full size (Greg Stark, Tom Lane)

The WAL update could be applied to the wrong page, potentially many pages past where it should have been. Aside from corrupting data, this error has been observed to result in significant “bloat” of standby servers compared to their masters, due to updates being applied far beyond where the end-of-file should have been. This failure mode does not appear to be a significant risk during crash recovery, only when initially synchronizing a standby created from a base backup taken from a quickly-changing master.

- Fix bug in determining when recovery has reached consistency (Tomonari Katsumata, Heikki Linnakangas)

In some cases WAL replay would mistakenly conclude that the database was already consistent at the start of replay, thus possibly allowing hot-standby queries before the database was really consistent. Other symptoms such as “PANIC: WAL contains references to invalid pages” were also possible.

- Fix improper locking of btree index pages while replaying a `VACUUM` operation in hot-standby mode (Andres Freund, Heikki Linnakangas, Tom Lane)

This error could result in “PANIC: WAL contains references to invalid pages” failures.

- Ensure that insertions into non-leaf GIN index pages write a full-page WAL record when appropriate (Heikki Linnakangas)

The previous coding risked index corruption in the event of a partial-page write during a system crash.
• When `pause_at_recovery_target` and `recovery_target_inclusive` are both set, ensure the target record is applied before pausing, not after (Heikki Linnakangas)

• Fix race conditions during server process exit (Robert Haas)

  Ensure that signal handlers don't attempt to use the process's MyProc pointer after it's no longer valid.

• Fix race conditions in walsender shutdown logic and walreceiver SIGHUP signal handler (Tom Lane)

• Fix unsafe references to `errno` within error reporting logic (Christian Kruse)

  This would typically lead to odd behaviors such as missing or inappropriate HINT fields.

• Fix possible crashes from using `ereport()` too early during server startup (Tom Lane)

  The principal case we've seen in the field is a crash if the server is started in a directory it doesn't have permission to read.

• Clear retry flags properly in OpenSSL socket write function (Alexander Kukushkin)

  This omission could result in a server lockup after unexpected loss of an SSL-encrypted connection.

• Fix length checking for Unicode identifiers (`U"..."` syntax) containing escapes (Tom Lane)

  A spurious truncation warning would be printed for such identifiers if the escaped form of the identifier was too long, but the identifier actually didn't need truncation after de-escaping.

• Allow keywords that are type names to be used in lists of roles (Stephen Frost)

  A previous patch allowed such keywords to be used without quoting in places such as role identifiers; but it missed cases where a list of role identifiers was permitted, such as `DROP ROLE`.

• Fix parser crash for `EXISTS(SELECT * FROM zero_column_table)` (Tom Lane)

• Fix possible crash due to invalid plan for nested sub-selects, such as `WHERE (... x IN (SELECT ...) ...) IN (SELECT ...) ` (Tom Lane)

• Fix `UPDATE/DELETE` of an inherited target table that has `UNION ALL` subqueries (Tom Lane)

  Without this fix, `UNION ALL` subqueries aren't correctly inserted into the update plans for inheritance child tables after the first one, typically resulting in no update happening for those child table(s).

• Ensure that `ANALYZE` creates statistics for a table column even when all the values in it are “too wide” (Tom Lane)

  `ANALYZE` intentionally omits very wide values from its histogram and most-common-values calculations, but it neglected to do something sane in the case that all the sampled entries are too wide.

• In `ALTER TABLE ... SET TABLESPACE`, allow the database's default tablespace to be used without a permissions check (Stephen Frost)

  `CREATE TABLE` has always allowed such usage, but `ALTER TABLE` didn't get the memo.

• Fix “cannot accept a set” error when some arms of a `CASE` return a set and others don't (Tom Lane)

• Properly distinguish numbers from non-numbers when generating JSON output (Andrew Dunstan)

• Fix checks for all-zero client addresses in `pgstat` functions (Kevin Grittner)

• Fix possible misclassification of multibyte characters by the text search parser (Tom Lane)
Non-ASCII characters could be misclassified when using C locale with a multibyte encoding. On Cygwin, non-C locales could fail as well.

- Fix possible misbehavior in `plainto_tsquery()` (Heikki Linnakangas)

  Use `memmove()` not `memcpy()` for copying overlapping memory regions. There have been no field reports of this actually causing trouble, but it's certainly risky.

- Fix placement of permissions checks in `pg_start_backup()` and `pg_stop_backup()` (Andres Freund, Magnus Hagander)

  The previous coding might attempt to do catalog access when it shouldn't.

- Accept `SHIFT_JIS` as an encoding name for locale checking purposes (Tatsuo Ishii)

- Fix `*`-qualification of named parameters in SQL-language functions (Tom Lane)

  Given a composite-type parameter named `foo`, `$1.*` worked fine, but `foo.*` not so much.

- Fix misbehavior of `PQhost()` on Windows (Fujii Masao)

  It should return `localhost` if no host has been specified.

- Improve error handling in `libpq` and `psql` for failures during `COPY TO STDOUT/FROM STDIN` (Tom Lane)

  In particular this fixes an infinite loop that could occur in 9.2 and up if the server connection was lost during `COPY FROM STDIN`. Variants of that scenario might be possible in older versions, or with other client applications.

- Fix incorrect translation handling in some `psql \d` commands (Peter Eisentraut, Tom Lane)

- Ensure `pg_basebackup`'s background process is killed when exiting its foreground process (Magnus Hagander)

- Fix possible incorrect printing of filenames in `pg_basebackup`'s verbose mode (Magnus Hagander)

- Avoid including tablespaces inside `PGDATA` twice in base backups (Dimitri Fontaine, Magnus Hagander)

- Fix misaligned descriptors in `ecpg` (MauMau)

- In `ecpg`, handle lack of a hostname in the connection parameters properly (Michael Meskes)

- Fix performance regression in `contrib/dblink` connection startup (Joe Conway)

  Avoid an unnecessary round trip when client and server encodings match.

- In `contrib/isn`, fix incorrect calculation of the check digit for ISMN values (Fabien Coelho)

- Fix `contrib/pg_stat_statement`'s handling of `CURRENT_DATE` and related constructs (Kyotaro Horiguchi)

- Ensure client-code-only installation procedure works as documented (Peter Eisentraut)

- In Mingw and Cygwin builds, install the `libpq DLL` in the `bin` directory (Andrew Dunstan)

  This duplicates what the MSVC build has long done. It should fix problems with programs like `psql` failing to start because they can't find the DLL.

- Avoid using the deprecated `dllwrap` tool in Cygwin builds (Marco Atzeri)
• Don’t generate plain-text HISTORY and src/test/regress/README files anymore (Tom Lane)

These text files duplicated the main HTML and PDF documentation formats. The trouble involved in maintaining them greatly outweighs the likely audience for plain-text format. Distribution tarballs will still contain files by these names, but they’ll just be stubs directing the reader to consult the main documentation. The plain-text INSTALL file will still be maintained, as there is arguably a use-case for that.

• Update time zone data files to tzdata release 2013i for DST law changes in Jordan and historical changes in Cuba.

In addition, the zones Asia/Riyadh87, Asia/Riyadh88, and Asia/Riyadh89 have been removed, as they are no longer maintained by IANA, and never represented actual civil timekeeping practice.

E.103. Release 9.2.6

Release date: 2013-12-05

This release contains a variety of fixes from 9.2.5. For information about new features in the 9.2 major release, see Section E.109.

E.103.1. Migration to Version 9.2.6

A dump/restore is not required for those running 9.2.X.

However, this release corrects a number of potential data corruption issues. See the first two changelog entries below to find out whether your installation has been affected and what steps you can take if so.

Also, if you are upgrading from a version earlier than 9.2.4, see Section E.105.

E.103.2. Changes

• Fix VACUUM’s tests to see whether it can update relfrozenxid (Andres Freund)

In some cases VACUUM (either manual or autovacuum) could incorrectly advance a table’s relfrozenxid value, allowing tuples to escape freezing, causing those rows to become invisible once 2^31 transactions have elapsed. The probability of data loss is fairly low since multiple incorrect advancements would need to happen before actual loss occurs, but it’s not zero. In 9.2.0 and later, the probability of loss is higher, and it’s also possible to get “could not access status of transaction” errors as a consequence of this bug. Users upgrading from releases 9.0.4 or 8.4.8 or earlier are not affected, but all later versions contain the bug.

The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix any latent corruption but will not be able to fix all pre-existing data errors. However, an installation can be presumed safe after performing this vacuuming if it has executed fewer than 2^31 update transactions in its lifetime (check this with SELECT txid_current() < 2^31).

• Fix initialization of pg_clog and pg_subtrans during hot standby startup (Andres Freund, Heikki Linnakangas)

This bug can cause data loss on standby servers at the moment they start to accept hot-standby queries, by marking committed transactions as uncommitted. The likelihood of such corruption is small unless, at the time of standby startup, the primary server has executed many updating transactions since its last checkpoint. Symptoms include missing rows, rows that should have been deleted being still visible, and obsolete versions of updated rows being still visible alongside their newer versions.
This bug was introduced in versions 9.3.0, 9.2.5, 9.1.10, and 9.0.14. Standby servers that have only been running earlier releases are not at risk. It's recommended that standby servers that have ever run any of the buggy releases be re- cloned from the primary (e.g., with a new base backup) after upgrading.

- Fix dangling-pointer problem in fast-path locking (Tom Lane)
  This could lead to corruption of the lock data structures in shared memory, causing “lock already held” and other odd errors.
- Truncate pg_multixact contents during WAL replay (Andres Freund)
  This avoids ever-increasing disk space consumption in standby servers.
- Ensure an anti-wraparound VACUUM counts a page as scanned when it's only verified that no tuples need freezing (Sergey Burladayan, Jeff Janes)
  This bug could result in failing to advance relfrozenxid, so that the table would still be thought to need another anti-wraparound vacuum. In the worst case the database might even shut down to prevent wraparound.
- Fix race condition in GIN index posting tree page deletion (Heikki Linnakangas)
  This could lead to transient wrong answers or query failures.
- Fix “unexpected spgdoinsert() failure” error during SP-GiST index creation (Teodor Sigaev)
- Avoid flattening a subquery whose SELECT list contains a volatile function wrapped inside a sub-SELECT (Tom Lane)
  This avoids unexpected results due to extra evaluations of the volatile function.
- Fix planner's processing of non-simple-variable subquery outputs nested within outer joins (Tom Lane)
  This error could lead to incorrect plans for queries involving multiple levels of subqueries within JOIN syntax.
- Fix incorrect planning in cases where the same non-strict expression appears in multiple WHERE and outer JOIN equality clauses (Tom Lane)
- Fix planner crash with whole-row reference to a subquery (Tom Lane)
- Fix incorrect generation of optimized MIN()/MAX() plans for inheritance trees (Tom Lane)
  The planner could fail in cases where the MIN()/MAX() argument was an expression rather than a simple variable.
- Fix premature deletion of temporary files (Andres Freund)
- Prevent intra-transaction memory leak when printing range values (Tom Lane)
  This fix actually cures transient memory leaks in any datatype output function, but range types are the only ones known to have had a significant problem.
- Prevent incorrect display of dropped columns in NOT NULL and CHECK constraint violation messages (Michael Paquier and Tom Lane)
- Allow default arguments and named-argument notation for window functions (Tom Lane)
  Previously, these cases were likely to crash.
• Fix possible read past end of memory in rule printing (Peter Eisentraut)
• Fix array slicing of `int2vector` and `oidvector` values (Tom Lane)

Expressions of this kind are now implicitly promoted to regular `int2` or `oid` arrays.
• Fix incorrect behaviors when using a SQL-standard, simple GMT offset timezone (Tom Lane)

In some cases, the system would use the simple GMT offset value when it should have used the regular timezone setting that had prevailed before the simple offset was selected. This change also causes the `timeofday` function to honor the simple GMT offset zone.
• Prevent possible misbehavior when logging translations of Windows error codes (Tom Lane)
• Properly quote generated command lines in `pg_ctl` (Naoya Anzai and Tom Lane)

This fix applies only to Windows.
• Fix `pg_dumpall` to work when a source database sets `default_transaction_read_only` via `ALTER DATABASE SET` (Kevin Grittner)

Previously, the generated script would fail during restore.
• Make `ecpg` search for quoted cursor names case-sensitively (Zoltán Böszörményi)
• Fix `ecpg`'s processing of lists of variables declared `varchar` (Zoltán Böszörményi)
• Make `contrib/lo` defend against incorrect trigger definitions (Marc Cousin)

Update time zone data files to tzdata release 2013h for DST law changes in Argentina, Brazil, Jordan, Libya, Liechtenstein, Morocco, and Palestine. Also, new timezone abbreviations WIB, WIT, WITA for Indonesia.

E.104. Release 9.2.5

Release date: 2013-10-10

This release contains a variety of fixes from 9.2.4. For information about new features in the 9.2 major release, see Section E.109.

E.104.1. Migration to Version 9.2.5

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.4, see Section E.105.

E.104.2. Changes

• Prevent corruption of multi-byte characters when attempting to case-fold identifiers (Andrew Dunstan)

PostgreSQL case-folds non-ASCII characters only when using a single-byte server encoding.
• Fix memory leak when creating B-tree indexes on range columns (Heikki Linnakangas)
• Fix checkpoint memory leak in background writer when `wal_level = hot_standby` (Naoya Anzai)
• Fix memory leak caused by `lo_open()` failure (Heikki Linnakangas)
• Fix memory overcommit bug when \texttt{work\_mem} is using more than 24GB of memory (Stephen Frost)

• Serializable snapshot fixes (Kevin Grittner, Heikki Linnakangas)

• Fix deadlock bug in libpq when using SSL (Stephen Frost)

• Fix possible SSL state corruption in threaded libpq applications (Nick Phillips, Stephen Frost)

• Improve estimate of planner cost when choosing between generic and custom plans (Tom Lane)
  This change will favor generic plans when planning cost is high.

• Properly compute row estimates for boolean columns containing many NULL values (Andrew Gierth)
  Previously tests like \texttt{col IS NOT TRUE} and \texttt{col IS NOT FALSE} did not properly factor in NULL values when estimating plan costs.

• Fix accounting for qualifier evaluation costs in UNION ALL and inheritance queries (Tom Lane)
  This fixes cases where suboptimal query plans could be chosen if some WHERE clauses are expensive to calculate.

• Prevent pushing down WHERE clauses into unsafe UNION/INTERSECT subqueries (Tom Lane)
  Subqueries of a UNION or INTERSECT that contain set-returning functions or volatile functions in their SELECT lists could be improperly optimized, leading to run-time errors or incorrect query results.

• Fix rare case of “failed to locate grouping columns” planner failure (Tom Lane)

• Fix \texttt{pg\_dump} of foreign tables with dropped columns (Andrew Dunstan)
  Previously such cases could cause a \texttt{pg\_upgrade} error.

• Reorder \texttt{pg\_dump} processing of extension-related rules and event triggers (Joe Conway)

• Force dumping of extension tables if specified by \texttt{pg\_dump -t} or \texttt{-n} (Joe Conway)

• Improve view dumping code's handling of dropped columns in referenced tables (Tom Lane)

• Fix \texttt{pg\_restore -l} with the directory archive to display the correct format name (Fujii Masao)

• Properly record index comments created using UNIQUE and PRIMARY KEY syntax (Andres Freund)
  This fixes a parallel \texttt{pg\_restore} failure.

• Cause \texttt{pg\_basebackup -x} with an empty xlog directory to throw an error rather than crashing (Magnus Hagander, Haruka Takatsuka)

• Properly guarantee transmission of WAL files before clean switchover (Fujii Masao)
  Previously, the streaming replication connection might close before all WAL files had been replayed on the standby.

• Fix WAL segment timeline handling during recovery (Mitsumasa Kondo, Heikki Linnakangas)
  WAL file recycling during standby recovery could lead to premature recovery completion, resulting in data loss.

• Prevent errors in WAL replay due to references to uninitialized empty pages (Andres Freund)
• Fix REINDEX TABLE and REINDEX DATABASE to properly revalidate constraints and mark invalidated indexes as valid (Noah Misch)

REINDEX INDEX has always worked properly.

• Avoid deadlocks during insertion into SP-GiST indexes (Teodor Sigaev)

• Fix possible deadlock during concurrent CREATE INDEX CONCURRENTLY operations (Tom Lane)

• Fix GiST index lookup crash (Tom Lane)

• Fix regexp_matches() handling of zero-length matches (Jeevan Chalke)

Previously, zero-length matches like "^" could return too many matches.

• Fix crash for overly-complex regular expressions (Heikki Linnakangas)

• Fix regular expression match failures for back references combined with non-greedy quantifiers (Jeevan Chalke)

• Prevent CREATE FUNCTION from checking SET variables unless function body checking is enabled (Tom Lane)

• Allow ALTER DEFAULT PRIVILEGES to operate on schemas without requiring CREATE permission (Tom Lane)

• Loosen restriction on keywords used in queries (Tom Lane)

Specifically, lessen keyword restrictions for role names, language names, EXPLAIN and COPY options, and SET values. This allows COPY ... (FORMAT BINARY) to work as expected; previously BINARY needed to be quoted.

• Print proper line number during COPY failure (Heikki Linnakangas)

• Fix pgp_pub_decrypt() so it works for secret keys with passwords (Marko Kreen)

• Make pg_upgrade use pg_dump --quote-all-identifiers to avoid problems with keyword changes between releases (Tom Lane)

• Remove rare inaccurate warning during vacuum of index-less tables (Heikki Linnakangas)

• Ensure that VACUUM ANALYZE still runs the ANALYZE phase if its attempt to truncate the file is cancelled due to lock conflicts (Kevin Grittner)

• Avoid possible failure when performing transaction control commands (e.g ROLLBACK) in prepared queries (Tom Lane)

• Ensure that floating-point data input accepts standard spellings of “infinity” on all platforms (Tom Lane)

The C99 standard says that allowable spellings are inf, +inf, -inf, infinity, +infinity, and -infinity. Make sure we recognize these even if the platform's strtod function doesn't.

• Avoid unnecessary reporting when track_activities is off (Tom Lane)

• Expand ability to compare rows to records and arrays (Rafal Rzepecki, Tom Lane)

• Prevent crash when psql's PSQLRC variable contains a tilde (Bruce Momjian)

• Add spinlock support for ARM64 (Mark Salter)

• Update time zone data files to tzdata release 2013d for DST law changes in Israel, Morocco, Palestine, and Paraguay. Also, historical zone data corrections for Macquarie Island.
E.105. Release 9.2.4

Release date: 2013-04-04

This release contains a variety of fixes from 9.2.3. For information about new features in the 9.2 major release, see Section E.109.

E.105.1. Migration to Version 9.2.4

A dump/restore is not required for those running 9.2.X.

However, this release corrects several errors in management of GiST indexes. After installing this update, it is advisable to \texttt{REINDEX} any GiST indexes that meet one or more of the conditions described below.

Also, if you are upgrading from a version earlier than 9.2.2, see Section E.107.

E.105.2. Changes

- Fix insecure parsing of server command-line switches (Mitsumasa Kondo, Kyotaro Horiguchi)

  A connection request containing a database name that begins with "--" could be crafted to damage or destroy files within the server's data directory, even if the request is eventually rejected. (CVE-2013-1899)

- Reset OpenSSL randomness state in each postmaster child process (Marko Kreen)

  This avoids a scenario wherein random numbers generated by \texttt{contrib/pgcrypto} functions might be relatively easy for another database user to guess. The risk is only significant when the postmaster is configured with \texttt{ssl=on} but most connections don't use SSL encryption. (CVE-2013-1900)

- Make \texttt{REPLICATION} privilege checks test current user not authenticated user (Noah Misch)

  An unprivileged database user could exploit this mistake to call \texttt{pg_start_backup()} or \texttt{pg_stop_backup()}, thus possibly interfering with creation of routine backups. (CVE-2013-1901)

- Fix GiST indexes to not use "fuzzy" geometric comparisons when it's not appropriate to do so (Alexander Korotkov)

  The core geometric types perform comparisons using "fuzzy" equality, but \texttt{gist_box_same} must do exact comparisons, else GiST indexes using it might become inconsistent. After installing this update, users should \texttt{REINDEX} any GiST indexes on \texttt{box}, \texttt{polygon}, \texttt{circle}, or \texttt{point} columns, since all of these use \texttt{gist_box_same}.

- Fix erroneous range-union and penalty logic in GiST indexes that use \texttt{contrib/btree_gist} for variable-width data types, that is \texttt{text}, \texttt{bytea}, \texttt{bit}, and \texttt{numeric} columns (Tom Lane)

  These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in useless index bloat. Users are advised to \texttt{REINDEX} such indexes after installing this update.

- Fix bugs in GiST page splitting code for multi-column indexes (Tom Lane)

  These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in indexes that are unnecessarily inefficient to search. Users are advised to \texttt{REINDEX} multi-column GiST indexes after installing this update.
• Fix `gist_point_consistent` to handle fuzziness consistently (Alexander Korotkov)

Index scans on GiST indexes on `point` columns would sometimes yield results different from a sequential scan, because `gist_point_consistent` disagreed with the underlying operator code about whether to do comparisons exactly or fuzzily.

• Fix buffer leak in WAL replay (Heikki Linnakangas)

This bug could result in “incorrect local pin count” errors during replay, making recovery impossible.

• Ensure we do crash recovery before entering archive recovery, if the database was not stopped cleanly and a `recovery.conf` file is present (Heikki Linnakangas, Kyotaro Horiguchi, Mitsumasa Kondo)

This is needed to ensure that the database is consistent in certain scenarios, such as initializing a standby server with a filesystem snapshot from a running server.

• Avoid deleting not-yet-archived WAL files during crash recovery (Heikki Linnakangas, Fujii Masao)

• Fix race condition in `DELETE RETURNING` (Tom Lane)

Under the right circumstances, `DELETE RETURNING` could attempt to fetch data from a shared buffer that the current process no longer has any pin on. If some other process changed the buffer meanwhile, this would lead to garbage `RETURNING` output, or even a crash.

• Fix infinite-loop risk in regular expression compilation (Tom Lane, Don Porter)

• Fix potential null-pointer dereference in regular expression compilation (Tom Lane)

• Fix `to_char()` to use ASCII-only case-folding rules where appropriate (Tom Lane)

This fixes misbehavior of some template patterns that should be locale-independent, but mishandled “I” and “i” in Turkish locales.

• Fix unwanted rejection of timestamp 1999-12-31 24:00:00 (Tom Lane)

• Fix SQL-language functions to be safely usable as support functions for range types (Tom Lane)

• Fix logic error when a single transaction does `UNLISTEN` then `LISTEN` (Tom Lane)

The session wound up not listening for notify events at all, though it surely should listen in this case.

• Fix possible planner crash after columns have been added to a view that's depended on by another view (Tom Lane)

• Fix performance issue in `EXPLAIN` (ANALYZE, TIMING OFF) (Pavel Stehule)

• Remove useless “picksplit doesn’t support secondary split” log messages (Josh Hansen, Tom Lane)

This message seems to have been added in expectation of code that was never written, and probably never will be, since GiST's default handling of secondary splits is actually pretty good. So stop nagging end users about it.

• Remove vestigial secondary-split support in `gist_box_picksplit()` (Tom Lane)

Not only was this implementation of secondary-split not better than the default implementation, it's actually worse. So remove it and let the default code path handle the case.

• Fix possible failure to send a session's last few transaction commit/abort counts to the statistics collector (Tom Lane)
• Eliminate memory leaks in PL/Perl's `spi_prepare()` function (Alex Hunsaker, Tom Lane)

• Fix `pg_dumpall` to handle database names containing “=” correctly (Heikki Linnakangas)

• Avoid crash in `pg_dump` when an incorrect connection string is given (Heikki Linnakangas)

• Ignore invalid indexes in `pg_dump` and `pg_upgrade` (Michael Paquier, Bruce Momjian)

  Dumping invalid indexes can cause problems at restore time, for example if the reason the index creation failed was because it tried to enforce a uniqueness condition not satisfied by the table's data. Also, if the index creation is in fact still in progress, it seems reasonable to consider it to be an uncommitted DDL change, which `pg_dump` wouldn't be expected to dump anyway. `pg_upgrade` now also skips invalid indexes rather than failing.

• In `pg_basebackup`, include only the current server version's subdirectory when backing up a table-space (Heikki Linnakangas)

• Add a server version check in `pg_basebackup` and `pg_receivexlog`, so they fail cleanly with version combinations that won't work (Heikki Linnakangas)

• Fix `contrib/dblink` to handle inconsistent settings of `DateStyle` or `IntervalStyle` safely (Daniel Farina, Tom Lane)

  Previously, if the remote server had different settings of these parameters, ambiguous dates might be read incorrectly. This fix ensures that datetime and interval columns fetched by a `dblink` query will be interpreted correctly. Note however that inconsistent settings are still risky, since literal values appearing in SQL commands sent to the remote server might be interpreted differently than they would be locally.

• Fix `contrib/pg_trgm`'s `similarity()` function to return zero for trigram-less strings (Tom Lane)

  Previously it returned NaN due to internal division by zero.

• Enable building PostgreSQL with Microsoft Visual Studio 2012 (Brar Piening, Noah Misch)

• Update time zone data files to `tzdata` release 2013b for DST law changes in Chile, Haiti, Morocco, Paraguay, and some Russian areas. Also, historical zone data corrections for numerous places.

  Also, update the time zone abbreviation files for recent changes in Russia and elsewhere: CHOT, GET, IRKT, KGT, KRAT, MAGT, MAWT, MSK, NOVT, OMST, TKT, VLAT, WST, YAKT, YEKT now follow their current meanings, and VOLT (Europe/Volgograd) and MIST (Antarctica/Macquarie) are added to the default abbreviations list.

### E.106. Release 9.2.3

**Release date:** 2013-02-07

This release contains a variety of fixes from 9.2.2. For information about new features in the 9.2 major release, see Section E.109.

### E.106.1. Migration to Version 9.2.3

A dump/restore is not required for those running 9.2.X.

However, if you are upgrading from a version earlier than 9.2.2, see Section E.107.

### E.106.2. Changes

• Prevent execution of `enum_recv` from SQL (Tom Lane)
The function was misdeclared, allowing a simple SQL command to crash the server. In principle an attacker might be able to use it to examine the contents of server memory. Our thanks to Sumit Soni (via Secunia SVCRP) for reporting this issue. (CVE-2013-0255)

- Fix multiple problems in detection of when a consistent database state has been reached during WAL replay (Fujii Masao, Heikki Linnakangas, Simon Riggs, Andres Freund)
- Fix detection of end-of-backup point when no actual redo work is required (Heikki Linnakangas)
  This mistake could result in incorrect “WAL ends before end of online backup” errors.
- Update minimum recovery point when truncating a relation file (Heikki Linnakangas)
  Once data has been discarded, it’s no longer safe to stop recovery at an earlier point in the timeline.
- Fix recycling of WAL segments after changing recovery target timeline (Heikki Linnakangas)
- Properly restore timeline history files from archive on cascading standby servers (Heikki Linnakangas)
- Fix lock conflict detection on hot-standby servers (Andres Freund, Robert Haas)
- Fix missing cancellations in hot standby mode (Noah Misch, Simon Riggs)
  The need to cancel conflicting hot-standby queries would sometimes be missed, allowing those queries to see inconsistent data.
- Prevent recovery pause feature from pausing before users can connect (Tom Lane)
- Fix SQL grammar to allow subscripting or field selection from a sub-SELECT result (Tom Lane)
- Fix performance problems with autovacuum truncation in busy workloads (Jan Wieck)

  Truncation of empty pages at the end of a table requires exclusive lock, but autovacuum was coded to fail (and release the table lock) when there are conflicting lock requests. Under load, it is easily possible that truncation would never occur, resulting in table bloat. Fix by performing a partial truncation, releasing the lock, then attempting to re-acquire the lock and continue. This fix also greatly reduces the average time before autovacuum releases the lock after a conflicting request arrives.
- Improve performance of SPI_execute and related functions, thereby improving PL/pgSQL's EXECUTE (Heikki Linnakangas, Tom Lane)

  Remove some data-copying overhead that was added in 9.2 as a consequence of revisions in the plan caching mechanism. This eliminates a performance regression compared to 9.1, and also saves memory, especially when the query string to be executed contains many SQL statements.

  A side benefit is that multi-statement query strings are now processed fully serially, that is we complete execution of earlier statements before running parse analysis and planning on the following ones. This eliminates a long-standing issue, in that DDL that should affect the behavior of a later statement will now behave as expected.
- Restore pre-9.2 cost estimates for index usage (Tom Lane)

  An ill-considered change of a fudge factor led to undesirably high cost estimates for use of very large indexes.
- Fix intermittent crash in DROP INDEX CONCURRENTLY (Tom Lane)
- Fix potential corruption of shared-memory lock table during CREATE/DROP INDEX CONCURRENTLY (Tom Lane)
• Fix COPY’s multiple-tuple-insertion code for the case of a tuple larger than page size minus fillfactor (Heikki Linnakangas)

The previous coding could get into an infinite loop.

• Protect against race conditions when scanning pg_tablespace (Stephen Frost, Tom Lane)

CREATE DATABASE and DROP DATABASE could misbehave if there were concurrent updates of pg_tablespace entries.

• Prevent DROP OWNED from trying to drop whole databases or tablespaces (Álvaro Herrera)

For safety, ownership of these objects must be reassigned, not dropped.

• Fix error in vacuum_freeze_table_age implementation (Andres Freund)

In installations that have existed for more than vacuum_freeze_min_age transactions, this mistake prevented autovacuum from using partial-table scans, so that a full-table scan would always happen instead.

• Prevent misbehavior when a RowExpr or XmlExpr is parse-analyzed twice (Andres Freund, Tom Lane)

This mistake could be user-visible in contexts such as CREATE TABLE LIKE INCLUDING INDEXES.

• Improve defenses against integer overflow in hashtable sizing calculations (Jeff Davis)

• Fix some bugs associated with privileges on datatypes (Tom Lane)

There were some issues with default privileges for types, and pg_dump failed to dump such privileges at all.

• Fix failure to ignore leftover temporary tables after a server crash (Tom Lane)

• Fix failure to rotate postmaster log files for size reasons on Windows (Jeff Janes, Heikki Linnakangas)

• Reject out-of-range dates in to_date() (Hitoshi Harada)

• Fix pg_extension_config_dump() to handle extension-update cases properly (Tom Lane)

This function will now replace any existing entry for the target table, making it usable in extension update scripts.

• Fix PL/pgSQL’s reporting of plan-time errors in possibly-simple expressions (Tom Lane)

The previous coding resulted in sometimes omitting the first line in the CONTEXT traceback for the error.

• Fix PL/Python’s handling of functions used as triggers on multiple tables (Andres Freund)

• Ensure that non-ASCII prompt strings are translated to the correct code page on Windows (Alexander Law, Noah Misch)

This bug affected psql and some other client programs.

• Fix possible crash in psql’s \? command when not connected to a database (Meng Qingzhong)

• Fix possible error if a relation file is removed while pg_basebackup is running (Heikki Linnakangas)

• Tolerate timeline switches while pg_basebackup -X fetch is backing up a standby server (Heikki Linnakangas)
• Make pg_dump exclude data of unlogged tables when running on a hot-standby server (Magnus Hagander)

This would fail anyway because the data is not available on the standby server, so it seems most convenient to assume --no-unlogged-table-data automatically.

• Fix pg_upgrade to deal with invalid indexes safely (Bruce Momjian)

• Fix pg_upgrade's -O/-o options (Marti Raudsepp)

• Fix one-byte buffer overrun in libpq's PQprintTuples (Xi Wang)

This ancient function is not used anywhere by PostgreSQL itself, but it might still be used by some client code.

• Make ecpglib use translated messages properly (Chen Huajun)

• Properly install ecpg_compat and pgtypes libraries on MSVC (Jiang Guiqing)

• Include our version of isinf() in libecpg if it's not provided by the system (Jiang Guiqing)

• Rearrange configure's tests for supplied functions so it is not fooled by bogus exports from libed-it/libreadline (Christoph Berg)

• Ensure Windows build number increases over time (Magnus Hagander)

• Make pgxs build executables with the right .exe suffix when cross-compiling for Windows (Zoltan Boszormenyi)

• Add new timezone abbreviation FET (Tom Lane)

This is now used in some eastern-European time zones.

E.107. Release 9.2.2

Release date: 2012-12-06

This release contains a variety of fixes from 9.2.1. For information about new features in the 9.2 major release, see Section E.109.

E.107.1. Migration to Version 9.2.2

A dump/restore is not required for those running 9.2.X.

However, you may need to perform REINDEX operations to correct problems in concurrently-built indexes, as described in the first changelog item below.

Also, if you are upgrading from version 9.2.0, see Section E.108.

E.107.2. Changes

• Fix multiple bugs associated with CREATE/DROP INDEX CONCURRENTLY (Andres Freund, Tom Lane, Simon Riggs, Pavan Deolasee)

An error introduced while adding DROP INDEX CONCURRENTLY allowed incorrect indexing decisions to be made during the initial phase of CREATE INDEX CONCURRENTLY; so that indexes built by that command could be corrupt. It is recommended that indexes built in 9.2.X with CREATE INDEX CONCURRENTLY be rebuilt after applying this update.

In addition, fix CREATE/DROP INDEX CONCURRENTLY to use in-place updates when changing the state of an index's pg_index row. This prevents race conditions that could cause concurrent
Release Notes

sessions to miss updating the target index, thus again resulting in corrupt concurrently-created indexes.

Also, fix various other operations to ensure that they ignore invalid indexes resulting from a failed CREATE INDEX CONCURRENTLY command. The most important of these is VACUUM, because an auto-vacuum could easily be launched on the table before corrective action can be taken to fix or remove the invalid index.

Also fix DROP INDEX CONCURRENTLY to not disable insertions into the target index until all queries using it are done.

Also fix misbehavior if DROP INDEX CONCURRENTLY is canceled: the previous coding could leave an un-droppable index behind.

• Correct predicate locking for DROP INDEX CONCURRENTLY (Kevin Grittner)

Previously, SSI predicate locks were processed at the wrong time, possibly leading to incorrect behavior of serializable transactions executing in parallel with the DROP.

• Fix buffer locking during WAL replay (Tom Lane)

The WAL replay code was insufficiently careful about locking buffers when replaying WAL records that affect more than one page. This could result in hot standby queries transiently seeing inconsistent states, resulting in wrong answers or unexpected failures.

• Fix an error in WAL generation logic for GIN indexes (Tom Lane)

This could result in index corruption, if a torn-page failure occurred.

• Fix an error in WAL replay logic for SP-GiST indexes (Tom Lane)

This could result in index corruption after a crash, or on a standby server.

• Fix incorrect detection of end-of-base-backup location during WAL recovery (Heikki Linnakangas)

This mistake allowed hot standby mode to start up before the database reaches a consistent state.

• Properly remove startup process's virtual XID lock when promoting a hot standby server to normal running (Simon Riggs)

This oversight could prevent subsequent execution of certain operations such as CREATE INDEX CONCURRENTLY.

• Avoid bogus “out-of-sequence timeline ID” errors in standby mode (Heikki Linnakangas)

• Prevent the postmaster from launching new child processes after it's received a shutdown signal (Tom Lane)

This mistake could result in shutdown taking longer than it should, or even never completing at all without additional user action.

• Fix the syslogger process to not fail when log_rotation_age exceeds 2^31 milliseconds (about 25 days) (Tom Lane)

• Fix WaitLatch() to return promptly when the requested timeout expires (Jeff Janes, Tom Lane)

With the previous coding, a steady stream of non-wait-terminating interrupts could delay return from WaitLatch() indefinitely. This has been shown to be a problem for the autovacuum launcher process, and might cause trouble elsewhere as well.

• Avoid corruption of internal hash tables when out of memory (Hitoshi Harada)
• Prevent file descriptors for dropped tables from being held open past transaction end (Tom Lane)

This should reduce problems with long-since-dropped tables continuing to occupy disk space.

• Prevent database-wide crash and restart when a new child process is unable to create a pipe for its latch (Tom Lane)

Although the new process must fail, there is no good reason to force a database-wide restart, so avoid that. This improves robustness when the kernel is nearly out of file descriptors.

• Avoid planner crash with joins to unflattened subqueries (Tom Lane)

• Fix planning of non-strict equivalence clauses above outer joins (Tom Lane)

The planner could derive incorrect constraints from a clause equating a non-strict construct to something else, for example WHERE COALESCE(foo, 0) = 0 when foo is coming from the nullable side of an outer join. 9.2 showed this type of error in more cases than previous releases, but the basic bug has been there for a long time.

• Fix SELECT DISTINCT with index-optimized MIN/MAX on an inheritance tree (Tom Lane)

The planner would fail with “failed to re-find MinMaxAggInfo record” given this combination of factors.

• Make sure the planner sees implicit and explicit casts as equivalent for all purposes, except in the minority of cases where there’s actually a semantic difference (Tom Lane)

• Include join clauses when considering whether partial indexes can be used for a query (Tom Lane)

A strict join clause can be sufficient to establish an x IS NOT NULL predicate, for example. This fixes a planner regression in 9.2, since previous versions could make comparable deductions.

• Limit growth of planning time when there are many indexable join clauses for the same index (Tom Lane)

• Improve planner’s ability to prove exclusion constraints from equivalence classes (Tom Lane)

• Fix partial-row matching in hashed subplans to handle cross-type cases correctly (Tom Lane)

This affects multicolumn NOT IN subplans, such as WHERE (a, b) NOT IN (SELECT x, y FROM ...) when for instance b and y are int4 and int8 respectively. This mistake led to wrong answers or crashes depending on the specific datatypes involved.

• Fix btree mark/restore functions to handle array keys (Tom Lane)

This oversight could result in wrong answers from merge joins whose inner side is an index scan using an indexed_column = ANY(array) condition.

• Revert patch for taking fewer snapshots (Tom Lane)

The 9.2 change to reduce the number of snapshots taken during query execution led to some anomalous behaviors not seen in previous releases, because execution would proceed with a snapshot acquired before locking the tables used by the query. Thus, for example, a query would not be guaranteed to see updates committed by a preceding transaction even if that transaction had exclusive lock. We’ll probably revisit this in future releases, but meanwhile put it back the way it was before 9.2.

• Acquire buffer lock when re-fetching the old tuple for an AFTER ROW UPDATE/DELETE trigger (Andres Freund)

In very unusual circumstances, this oversight could result in passing incorrect data to a trigger WHEN condition, or to the precheck logic for a foreign-key enforcement trigger. That could result in a crash, or in an incorrect decision about whether to fire the trigger.
• Fix ALTER COLUMN TYPE to handle inherited check constraints properly (Pavan Deolasee)
  This worked correctly in pre-8.4 releases, and now works correctly in 8.4 and later.

• Fix ALTER EXTENSION SET SCHEMA’s failure to move some subsidiary objects into the new
  schema (Álvaro Herrera, Dimitri Fontaine)

• Handle CREATE TABLE AS EXECUTE correctly in extended query protocol (Tom Lane)

• Don’t modify the input parse tree in DROP RULE IF NOT EXISTS and DROP TRIGGER IF
  NOT EXISTS (Tom Lane)
  This mistake would cause errors if a cached statement of one of these types was re-executed.

• Fix REASSIGN OWNED to handle grants on tablespaces (Álvaro Herrera)

• Ignore incorrect pg_attribute entries for system columns for views (Tom Lane)
  Views do not have any system columns. However, we forgot to remove such entries when converting
  a table to a view. That’s fixed properly for 9.3 and later, but in previous branches we need to defend
  against existing mis-converted views.

• Fix rule printing to dump INSERT INTO table DEFAULT VALUES correctly (Tom Lane)

• Guard against stack overflow when there are too many UNION/INTERSECT/EXCEPT clauses in
  a query (Tom Lane)

• Prevent platform-dependent failures when dividing the minimum possible integer value by -1 (Xi
  Wang, Tom Lane)

• Fix possible access past end of string in date parsing (Hitoshi Harada)

• Fix failure to advance XID epoch if XID wraparound happens during a checkpoint and wal_level
  is hot_standby (Tom Lane, Andres Freund)
  While this mistake had no particular impact on PostgreSQL itself, it was bad for applications that
  rely on txid_current() and related functions: the TXID value would appear to go backwards.

• Fix pg_terminate_backend() and pg_cancel_backend() to not throw error for a non-
  existent target process (Josh Kupershmidt)
  This case already worked as intended when called by a superuser, but not so much when called by
  ordinary users.

• Fix display of pg_stat_replication.sync_state at a page boundary (Kyotaro Horiguchi)

• Produce an understandable error message if the length of the path name for a Unix-domain socket
  exceeds the platform-specific limit (Tom Lane, Andrew Dunstan)
  Formerly, this would result in something quite unhelpful, such as “Non-recoverable failure in name
  resolution”.

• Fix memory leaks when sending composite column values to the client (Tom Lane)

• Save some cycles by not searching for subtransaction locks at commit (Simon Riggs)
  In a transaction holding many exclusive locks, this useless activity could be quite costly.

• Make pg_ctl more robust about reading the postmaster.pid file (Heikki Linnakangas)
  This fixes race conditions and possible file descriptor leakage.
• Fix possible crash in psql if incorrectly-encoded data is presented and the client_encoding setting is a client-only encoding, such as SJIS (Jiang Guiqing)

• Make pg_dump dump SEQUENCE SET items in the data not pre-data section of the archive (Tom Lane)

This fixes an undesirable inconsistency between the meanings of --data-only and --section=data, and also fixes dumping of sequences that are marked as extension configuration tables.

• Fix pg_dump's handling of DROP DATABASE commands in --clean mode (Guillaume Lelarge)

Beginning in 9.2.0, pg_dump --clean would issue a DROP DATABASE command, which was either useless or dangerous depending on the usage scenario. It no longer does that. This change also fixes the combination of --clean and --create to work sensibly, i.e., emit DROP DATABASE then CREATE DATABASE before reconnecting to the target database.

• Fix pg_dump for views with circular dependencies and no relation options (Tom Lane)

The previous fix to dump relation options when a view is involved in a circular dependency didn’t work right for the case that the view has no options; it emitted ALTER VIEW foo SET () which is invalid syntax.

• Fix bugs in the restore.sql script emitted by pg_dump in tar output format (Tom Lane)

The script would fail outright on tables whose names include upper-case characters. Also, make the script capable of restoring data in --inserts mode as well as the regular COPY mode.

• Fix pg_dump to accept POSIX-conformant tar files (Brian Weaver, Tom Lane)

The original coding of pg_dump's tar output mode produced files that are not fully conformant with the POSIX standard. This has been corrected for version 9.3. This patch updates previous branches so that they will accept both the incorrect and the corrected formats, in hopes of avoiding compatibility problems when 9.3 comes out.

• Fix tar files emitted by pg_basebackup to be POSIX conformant (Brian Weaver, Tom Lane)

• Fix pg_resetxlog to locate postmaster.pid correctly when given a relative path to the data directory (Tom Lane)

This mistake could lead to pg_resetxlog not noticing that there is an active postmaster using the data directory.

• Fix libpq's lo_import() and lo_export() functions to report file I/O errors properly (Tom Lane)

• Fix ecpg's processing of nested structure pointer variables (Muhammad Usama)

• Fix ecpg's ecpg_get_data function to handle arrays properly (Michael Meskes)

• Prevent pg_upgrade from trying to process TOAST tables for system catalogs (Bruce Momjian)

This fixes an error seen when the information_schema has been dropped and recreated. Other failures were also possible.

• Improve pg_upgrade performance by setting synchronous_commit to off in the new cluster (Bruce Momjian)

• Make contrib/pageinspect's btree page inspection functions take buffer locks while examining pages (Tom Lane)

• Work around unportable behavior of malloc(0) and realloc(NULL, 0) (Tom Lane)
On platforms where these calls return NULL, some code mistakenly thought that meant out-of-memory. This is known to have broken pg_dump for databases containing no user-defined aggregates. There might be other cases as well.

- Ensure that make install for an extension creates the extension installation directory (Cédric Villemain)
  
  Previously, this step was missed if MODuledir was set in the extension's Makefile.

- Fix pgxs support for building loadable modules on AIX (Tom Lane)
  
  Building modules outside the original source tree didn't work on AIX.

- Update time zone data files to tzdata release 2012j for DST law changes in Cuba, Israel, Jordan, Libya, Palestine, Western Samoa, and portions of Brazil.

**E.108. Release 9.2.1**

**Release date:** 2012-09-24

This release contains a variety of fixes from 9.2.0. For information about new features in the 9.2 major release, see Section E.109.

**E.108.1. Migration to Version 9.2.1**

A dump/restore is not required for those running 9.2.X.

However, you may need to perform REINDEX and/or VACUUM operations to recover from the effects of the data corruption bug described in the first changelog item below.

**E.108.2. Changes**

- Fix persistence marking of shared buffers during WAL replay (Jeff Davis)
  
  This mistake can result in buffers not being written out during checkpoints, resulting in data corruption if the server later crashes without ever having written those buffers. Corruption can occur on any server following crash recovery, but it is significantly more likely to occur on standby slave servers since those perform much more WAL replay. There is a low probability of corruption of btree and GIN indexes. There is a much higher probability of corruption of table “visibility maps”, which might lead to wrong answers from index-only scans. Table data proper cannot be corrupted by this bug.

  While no index corruption due to this bug is known to have occurred in the field, as a precautionary measure it is recommended that production installations REINDEX all btree and GIN indexes at a convenient time after upgrading to 9.2.1.

  Also, it is recommended to perform a VACUUM of all tables while having vacuum_freeze_table_age set to zero. This will fix any incorrect visibility map data. vacuum_cost_delay can be adjusted to reduce the performance impact of vacuuming, while causing it to take longer to finish.

- Fix possible incorrect sorting of output from queries involving WHERE indexed_column IN (list_of_values) (Tom Lane)

- Fix planner failure for queries involving GROUP BY expressions along with window functions and aggregates (Tom Lane)

- Fix planner's assignment of executor parameters (Tom Lane)
This error could result in wrong answers from queries that scan the same WITH subquery multiple times.

- Improve planner's handling of join conditions in index scans (Tom Lane)
- Improve selectivity estimation for text search queries involving prefixes, i.e. \textit{word:*} patterns (Tom Lane)
- Fix delayed recognition of permissions changes (Tom Lane)

A command that needed no locks other than ones its transaction already had might fail to notice a concurrent \texttt{GRANT} or \texttt{REVOKE} that committed since the start of its transaction.

- Fix \texttt{ANALYZE} to not fail when a column is a domain over an array type (Tom Lane)
- Prevent PL/Perl from crashing if a recursive PL/Perl function is redefined while being executed (Tom Lane)
- Work around possible misoptimization in PL/Perl (Tom Lane)

Some Linux distributions contain an incorrect version of \texttt{pthread.h} that results in incorrect compiled code in PL/Perl, leading to crashes if a PL/Perl function calls another one that throws an error.

- Remove unnecessary dependency on \texttt{pg_config} from \texttt{pg_upgrade} (Peter Eisentraut)
- Update time zone data files to \texttt{tzdata} release 2012f for DST law changes in Fiji

\section*{E.109. Release 9.2}

\textbf{Release date:} 2012-09-10

\subsection*{E.109.1. Overview}

This release has been largely focused on performance improvements, though new SQL features are not lacking. Work also continues in the area of replication support. Major enhancements include:

- Allow queries to retrieve data only from indexes, avoiding heap access (\textit{index-only scans})
- Allow the planner to generate custom plans for specific parameter values even when using prepared statements
- Improve the planner's ability to use nested loops with inner index scans
- Allow streaming replication slaves to forward data to other slaves (\textit{cascading replication})
- Allow \texttt{pg_basebackup} to make base backups from standby servers
- Add a \texttt{pg_receivexlog} tool to archive WAL file changes as they are written
- Add the SP-GiST (Space-Partitioned GiST) index access method
- Add support for range data types
- Add a \texttt{JSON} data type
- Add a \texttt{security_barrier} option for views
- Allow \texttt{libpq} connection strings to have the format of a URI
- Add a single-row processing mode to \texttt{libpq} for better handling of large result sets

The above items are explained in more detail in the sections below.
E.109.2. Migration to Version 9.2

A dump/restore using pg_dump, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.2 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

E.109.2.1. System Catalogs

- Remove the spclocation field from pg_tablespace (Magnus Hagander)

  This field was duplicative of the symbolic links that actually define tablespace locations, and thus risked errors of omission when moving a tablespace. This change allows tablespace directories to be moved while the server is down, by manually adjusting the symbolic links. To replace this field, we have added pg_tablespace_location() to allow querying of the symbolic links.

- Move tsvector most-common-element statistics to new pg_stats columns (Alexander Kordotkov)

  Consult most_common_elements and most_common_elem_freqs for the data formerly available in most_common_vals and most_common_freqs for a tsvector column.

E.109.2.2. Functions

- Remove hstore's => operator (Robert Haas)

  Users should now use hstore(text, text). Since PostgreSQL 9.0, a warning message has been emitted when an operator named => is created because the SQL standard reserves that token for another use.

- Ensure that xpath() escapes special characters in string values (Florian Pflug)

  Without this it is possible for the result not to be valid XML.

- Make pg_relation_size() and friends return NULL if the object does not exist (Phil Sorber)

  This prevents queries that call these functions from returning errors immediately after a concurrent DROP.

- Make EXTRACT(EPOCH FROM timestamp without time zone) measure the epoch from local midnight, not UTC midnight (Tom Lane)

  This change reverts an ill-considered change made in release 7.3. Measuring from UTC midnight was inconsistent because it made the result dependent on the timezone setting, which computations for timestamp without time zone should not be. The previous behavior remains available by casting the input value to timestamp with time zone.

- Properly parse time strings with trailing yesterday, today, and tomorrow (Dean Rasheed)

  Previously, SELECT '04:00:00 yesterday'::timestamp returned yesterday's date at midnight.

- Fix to_date() and to_timestamp() to wrap incomplete dates toward 2020 (Bruce Momijian)

  Previously, supplied years and year masks of less than four digits wrapped inconsistently.

E.109.2.3. Object Modification

- Prevent ALTER DOMAIN from working on non-domain types (Peter Eisentraut)
Owner and schema changes were previously possible on non-domain types.

- No longer forcibly lowercase procedural language names in `CREATE FUNCTION` (Robert Haas)

  While unquoted language identifiers are still lowercased, strings and quoted identifiers are no longer forcibly down-cased. Thus for example `CREATE FUNCTION ... LANGUAGE 'C'` will no longer work; it must be spelled `'c'`, or better omit the quotes.

- Change system-generated names of foreign key enforcement triggers (Tom Lane)

  This change ensures that the triggers fire in the correct order in some corner cases involving self-referential foreign key constraints.

**E.109.2.4. Command-Line Tools**

- Provide consistent backquote, variable expansion, and quoted substring behavior in `psql` meta-command arguments (Tom Lane)

  Previously, such references were treated oddly when not separated by whitespace from adjacent text. For example `'FOO'BAR` was output as `FOO BAR` (unexpected insertion of a space) and `FOO'BAR'BAZ` was output unchanged (not removing the quotes as most would expect).

- No longer treat `clusterdb` table names as double-quoted; no longer treat `reindexdb` table and index names as double-quoted (Bruce Momjian)

  Users must now include double-quotes in the command arguments if quoting is wanted.

- `createuser` no longer prompts for option settings by default (Peter Eisentraut)

  Use `--interactive` to obtain the old behavior.

- Disable prompting for the user name in `dropuser` unless `--interactive` is specified (Peter Eisentraut)

**E.109.2.5. Server Settings**

- Add server parameters for specifying the locations of server-side SSL files (Peter Eisentraut)

  This allows changing the names and locations of the files that were previously hard-coded as `server.crt`, `server.key`, `root.crt`, and `root.crl` in the data directory. *The server will no longer examine `root.crt` or `root.crl` by default*; to load these files, the associated parameters must be set to non-default values.

- Remove the `silent_mode` parameter (Heikki Linnakangas)

  Similar behavior can be obtained with `pg_ctl start -l postmaster.log`.

- Remove the `wal_sender_delay` parameter, as it is no longer needed (Tom Lane)

- Remove the `custom_variable_classes` parameter (Tom Lane)

  The checking provided by this setting was dubious. Now any setting can be prefixed by any class name.

**E.109.2.6. Monitoring**

- Rename `pg_stat_activity.procpid` to `pid`, to match other system tables (Magnus Hagander)

- Create a separate `pg_stat_activity` column to report process state (Scott Mead, Magnus Hagander)
The previous query and query_start values now remain available for an idle session, allowing enhanced analysis.

- Rename pg_stat_activity.current_query to query because it is not cleared when the query completes (Magnus Hagander)

- Change all SQL-level statistics timing values to be float8 columns measured in milliseconds (Tom Lane)

This change eliminates the designed-in assumption that the values are accurate to microseconds and no more (since the float8 values can be fractional). The columns affected are pg_stat_user_functions.total_time, pg_stat_user_functions.self_time, pg_stat_xact_user_functions.total_time, and pg_stat_xact_user_functions.self_time. The statistics functions underlying these columns now also return float8 milliseconds, rather than bigint microseconds. contrib/pg_stat_statements' total_time column is now also measured in milliseconds.

E.109.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.2 and the previous major release.

E.109.3.1. Server

E.109.3.1.1. Performance

- Allow queries to retrieve data only from indexes, avoiding heap access (Robert Haas, Ibrar Ahmed, Heikki Linnakangas, Tom Lane)

  This feature is often called index-only scans. Heap access can be skipped for heap pages containing only tuples that are visible to all sessions, as reported by the visibility map; so the benefit applies mainly to mostly-static data. The visibility map was made crash-safe as a necessary part of implementing this feature.

- Add the SP-GiST (Space-Partitioned GiST) index access method (Teodor Sigaev, Oleg Bartunov, Tom Lane)

  SP-GiST is comparable to GiST in flexibility, but supports unbalanced partitioned search structures rather than balanced trees. For suitable problems, SP-GiST can be faster than GiST in both index build time and search time.

- Allow group commit to work effectively under heavy load (Peter Geoghegan, Simon Riggs, Heikki Linnakangas)

  Previously, batching of commits became ineffective as the write workload increased, because of internal lock contention.

- Allow uncontended locks to be managed using a new fast-path lock mechanism (Robert Haas)

- Reduce overhead of creating virtual transaction ID locks (Robert Haas)

- Reduce the overhead of serializable isolation level locks (Dan Ports)

- Improve PowerPC and Itanium spinlock performance (Manabu Ori, Robert Haas, Tom Lane)

- Reduce overhead for shared invalidation cache messages (Robert Haas)

- Move the frequently accessed members of the PGPROC shared memory array to a separate array (Pavan Deolasee, Heikki Linnakangas, Robert Haas)
• Improve `COPY` performance by adding tuples to the heap in batches (Heikki Linnakangas)

• Improve GiST index performance for geometric data types by producing better trees with less memory allocation overhead (Alexander Korotkov)

• Improve GiST index build times (Alexander Korotkov, Heikki Linnakangas)

• Allow hint bits to be set sooner for temporary and unlogged tables (Robert Haas)

• Allow sorting to be performed by inlined, non-SQL-callable comparison functions (Peter Geoghegan, Robert Haas, Tom Lane)

• Make the number of CLOG buffers scale based on `shared_buffers` (Robert Haas, Simon Riggs, Tom Lane)

• Improve performance of buffer pool scans that occur when tables or databases are dropped (Jeff Janes, Simon Riggs)

• Improve performance of checkpointer’s fsync-request queue when many tables are being dropped or truncated (Tom Lane)

• Pass the safe number of file descriptors to child processes on Windows (Heikki Linnakangas)

This allows Windows sessions to use more open file descriptors than before.

**E.109.3.1.2. Process Management**

• Create a dedicated background process to perform checkpoints (Simon Riggs)

Formerly the background writer did both dirty-page writing and checkpointing. Separating this into two processes allows each goal to be accomplished more predictably.

• Improve asynchronous commit behavior by waking the walwriter sooner (Simon Riggs)

Previously, only `wal_writer_delay` triggered WAL flushing to disk; now filling a WAL buffer also triggers WAL writes.

• Allow the bgwriter, walwriter, checkpointer, statistics collector, log collector, and archiver background processes to sleep more efficiently during periods of inactivity (Peter Geoghegan, Tom Lane)

This series of changes reduces the frequency of process wake-ups when there is nothing to do, dramatically reducing power consumption on idle servers.

**E.109.3.1.3. Optimizer**

• Allow the planner to generate custom plans for specific parameter values even when using prepared statements (Tom Lane)

In the past, a prepared statement always had a single “generic” plan that was used for all parameter values, which was frequently much inferior to the plans used for non-prepared statements containing explicit constant values. Now, the planner attempts to generate custom plans for specific parameter values. A generic plan will only be used after custom plans have repeatedly proven to provide no benefit. This change should eliminate the performance penalties formerly seen from use of prepared statements (including non-dynamic statements in PL/pgSQL).

• Improve the planner’s ability to use nested loops with inner index scans (Tom Lane)

The new “parameterized path” mechanism allows inner index scans to use values from relations that are more than one join level up from the scan. This can greatly improve performance in situations where semantic restrictions (such as outer joins) limit the allowed join orderings.
• Improve the planning API for foreign data wrappers (Etsuro Fujita, Shigeru Hanada, Tom Lane)
  Wrappers can now provide multiple access “paths” for their tables, allowing more flexibility in join planning.

• Recognize self-contradictory restriction clauses for non-table relations (Tom Lane)
  This check is only performed when constraint_exclusion is on.

• Allow indexed_col op ANY(ARRAY[...]) conditions to be used in plain index scans and index-only scans (Tom Lane)
  Formerly such conditions could only be used in bitmap index scans.

• Support MIN/MAX index optimizations on boolean columns (Marti Raudsepp)

• Account for set-returning functions in SELECT target lists when setting row count estimates (Tom Lane)

• Fix planner to handle indexes with duplicated columns more reliably (Tom Lane)

• Collect and use element-frequency statistics for arrays (Alexander Korotkov, Tom Lane)
  This change improves selectivity estimation for the array <@, &&, and @> operators (array containment and overlaps).

• Allow statistics to be collected for foreign tables (Etsuro Fujita)

• Improve cost estimates for use of partial indexes (Tom Lane)

• Improve the planner's ability to use statistics for columns referenced in subqueries (Tom Lane)

• Improve statistical estimates for subqueries using DISTINCT (Tom Lane)

E.109.3.1.4. Authentication

• Do not treat role names and samerole specified in pg_hba.conf as automatically including superusers (Andrew Dunstan)
  This makes it easier to use reject lines with group roles.

• Adjust pg_hba.conf processing to handle token parsing more consistently (Brendan Jurd, Álvaro Herrera)

•Disallow empty pg_hba.conf files (Tom Lane)
  This was done to more quickly detect misconfiguration.

• Make superuser privilege imply replication privilege (Noah Misch)
  This avoids the need to explicitly assign such privileges.

E.109.3.1.5. Monitoring

• Attempt to log the current query string during a backend crash (Marti Raudsepp)
• Make logging of autovacuum I/O activity more verbose (Greg Smith, Noah Misch)
  This logging is triggered by log_autovacuum_min_duration.

• Make WAL replay report failures sooner (Fujii Masao)
  There were some cases where failures were only reported once the server went into master mode.
• Add `pg_xlog_location_diff()` to simplify WAL location comparisons (Euler Taveira de Oliveira)

This is useful for computing replication lag.

• Support configurable event log application names on Windows (MauMau, Magnus Hagander)

This allows different instances to use the event log with different identifiers, by setting the `event_source` server parameter, which is similar to how `syslog_ident` works.

• Change “unexpected EOF” messages to `DEBUG1` level, except when there is an open transaction (Magnus Hagander)

This change reduces log chatter caused by applications that close database connections ungracefully.

E.109.3.1.6. Statistical Views

• Track temporary file sizes and file counts in the `pg_stat_database` system view (Tomas Von dra)

• Add a deadlock counter to the `pg_stat_database` system view (Magnus Hagander)

• Add a server parameter `track_io_timing` to track I/O timings (Ants Aasma, Robert Haas)

• Report checkpoint timing information in `pg_stat_bgwriter` (Greg Smith, Peter Geoghegan)

E.109.3.1.7. Server Settings

• Silently ignore nonexistent schemas specified in `search_path` (Tom Lane)

This makes it more convenient to use generic path settings, which might include some schemas that don’t exist in all databases.

• Allow superusers to set `deadlock_timeout` per-session, not just per-cluster (Noah Misch)

This allows `deadlock_timeout` to be reduced for transactions that are likely to be involved in a deadlock, thus detecting the failure more quickly. Alternatively, increasing the value can be used to reduce the chances of a session being chosen for cancellation due to a deadlock.

• Add a server parameter `temp_file_limit` to constrain temporary file space usage per session (Mark Kirkwood)

• Allow a superuser to `SET` an extension’s superuser-only custom variable before loading the associated extension (Tom Lane)

The system now remembers whether a `SET` was performed by a superuser, so that proper privilege checking can be done when the extension is loaded.

• Add postmaster `-C` option to query configuration parameters (Bruce Momjian)

This allows `pg_ctl` to better handle cases where `PGDATA` or `-D` points to a configuration-only directory.

• Replace an empty locale name with the implied value in `CREATE DATABASE` (Tom Lane)

This prevents cases where `pg_database.datcollate` or `datctype` could be interpreted differently after a server restart.

E.109.3.1.7.1. `postgresql.conf`

• Allow multiple errors in `postgresql.conf` to be reported, rather than just the first one (Alexey Klyukin, Tom Lane)
• Allow a reload of `postgresql.conf` to be processed by all sessions, even if there are some settings that are invalid for particular sessions (Alexey Klyukin)

Previously, such not-valid-within-session values would cause all setting changes to be ignored by that session.

• Add an `include_if_exists` facility for configuration files (Greg Smith)

This works the same as `include`, except that an error is not thrown if the file is missing.

• Identify the server time zone during `initdb`, and set `postgresql.conf` entries `timezone` and `log_timezone` accordingly (Tom Lane)

This avoids expensive time zone probes during server start.

• Fix `pg_settings` to report `postgresql.conf` line numbers on Windows (Tom Lane)

E.109.3.2. Replication and Recovery

• Allow streaming replication slaves to forward data to other slaves (`cascading replication`) (Fujii Masao)

Previously, only the master server could supply streaming replication log files to standby servers.

• Add new `synchronous_commit` mode `remote_write` (Fujii Masao, Simon Riggs)

This mode waits for the standby server to write transaction data to its own operating system, but does not wait for the data to be flushed to the standby's disk.

• Add a `pg_receivexlog` tool to archive WAL file changes as they are written, rather than waiting for completed WAL files (Magnus Hagander)

• Allow `pg_basebackup` to make base backups from standby servers (Jun Ishizuka, Fujii Masao)

This feature lets the work of making new base backups be off-loaded from the primary server.

• Allow streaming of WAL files while `pg_basebackup` is performing a backup (Magnus Hagander)

This allows passing of WAL files to the standby before they are discarded on the primary.

E.109.3.3. Queries

• Cancel the running query if the client gets disconnected (Florian Pflug)

If the backend detects loss of client connection during a query, it will now cancel the query rather than attempting to finish it.

• Retain column names at run time for row expressions (Andrew Dunstan, Tom Lane)

This change allows better results when a row value is converted to `hstore` or `json` type: the fields of the resulting value will now have the expected names.

• Improve column labels used for sub-`SELECT` results (Marti Raudsepp)

Previously, the generic label `?column?` was used.

• Improve heuristics for determining the types of unknown values (Tom Lane)

The longstanding rule that an unknown constant might have the same type as the value on the other side of the operator using it is now applied when considering polymorphic operators, not only for simple operator matches.
• Warn about creating casts to or from domain types (Robert Haas)

Such casts have no effect.

• When a row fails a CHECK or NOT NULL constraint, show the row’s contents as error detail (Jan Kundrát)

This should make it easier to identify which row is problematic when an insert or update is processing many rows.

E.109.3.4. Object Manipulation

• Provide more reliable operation during concurrent DDL (Robert Haas, Noah Misch)

This change adds locking that should eliminate “cache lookup failed” errors in many scenarios. Also, it is no longer possible to add relations to a schema that is being concurrently dropped, a scenario that formerly led to inconsistent system catalog contents.

• Add CONCURRENTLY option to DROP INDEX (Simon Riggs)

This allows index removal without blocking other sessions.

• Allow foreign data wrappers to have per-column options (Shigeru Hanada)

• Improve pretty-printing of view definitions (Andrew Dunstan)

E.109.3.4.1. Constraints

• Allow CHECK constraints to be declared NOT VALID (Álvaro Herrera)

Adding a NOT VALID constraint does not cause the table to be scanned to verify that existing rows meet the constraint. Subsequently, newly added or updated rows are checked. Such constraints are ignored by the planner when considering constraint_exclusion, since it is not certain that all rows meet the constraint.

The new ALTER TABLE VALIDATE command allows NOT VALID constraints to be checked for existing rows, after which they are converted into ordinary constraints.

• Allow CHECK constraints to be declared NO INHERIT (Nikhil Sontakke, Alex Hunsaker, Álvaro Herrera)

This makes them enforceable only on the parent table, not on child tables.

• Add the ability to rename constraints (Peter Eisentraut)

E.109.3.4.2. ALTER

• Reduce need to rebuild tables and indexes for certain ALTER TABLE ... ALTER COLUMN TYPE operations (Noah Misch)

Increasing the length limit for a varchar or varbit column, or removing the limit altogether, no longer requires a table rewrite. Similarly, increasing the allowable precision of a numeric column, or changing a column from constrained numeric to unconstrained numeric, no longer requires a table rewrite. Table rewrites are also avoided in similar cases involving the interval, timestamp, and timestamptz types.

• Avoid having ALTER TABLE revalidate foreign key constraints in some cases where it is not necessary (Noah Misch)

• Add IF EXISTS options to some ALTER commands (Pavel Stehule)

For example, ALTER FOREIGN TABLE IF EXISTS foo RENAME TO bar.
• Add ALTER FOREIGN DATA WRAPPER ... RENAME and ALTER SERVER ... RENAME (Peter Eisentraut)

• Add ALTER DOMAIN ... RENAME (Peter Eisentraut)

You could already rename domains using ALTER TYPE.

• Throw an error for ALTER DOMAIN ... DROP CONSTRAINT on a nonexistent constraint (Peter Eisentraut)

An IF EXISTS option has been added to provide the previous behavior.

E.109.3.4.3. CREATE TABLE

• Allow CREATE TABLE (LIKE ...) from foreign tables, views, and composite types (Peter Eisentraut)

For example, this allows a table to be created whose schema matches a view.

• Fix CREATE TABLE (LIKE ...) to avoid index name conflicts when copying index comments (Tom Lane)

• Fix CREATE TABLE ... AS EXECUTE to handle WITH NO DATA and column name specifications (Tom Lane)

E.109.3.4.4. Object Permissions

• Add a security_barrier option for views (KaiGai Kohei, Robert Haas)

This option prevents optimizations that might allow view-protected data to be exposed to users, for example pushing a clause involving an insecure function into the WHERE clause of the view. Such views can be expected to perform more poorly than ordinary views.

• Add a new LEAKPROOF function attribute to mark functions that can safely be pushed down into security_barrier views (KaiGai Kohei)

• Add support for privileges on data types (Peter Eisentraut)

This adds support for the SQL-conforming USAGE privilege on types and domains. The intent is to be able to restrict which users can create dependencies on types, since such dependencies limit the owner's ability to alter the type.

• Check for INSERT privileges in SELECT INTO / CREATE TABLE AS (KaiGai Kohei)

Because the object is being created by SELECT INTO or CREATE TABLE AS, the creator would ordinarily have insert permissions; but there are corner cases where this is not true, such as when ALTER DEFAULT PRIVILEGES has removed such permissions.

E.109.3.5. Utility Operations

• Allow VACUUM to more easily skip pages that cannot be locked (Simon Riggs, Robert Haas)

This change should greatly reduce the incidence of VACUUM getting “stuck” waiting for other sessions.

• Make EXPLAIN (BUFFERS) count blocks dirtied and written (Robert Haas)

• Make EXPLAIN ANALYZE report the number of rows rejected by filter steps (Marko Tiikkaja)

• Allow EXPLAIN ANALYZE to avoid timing overhead when time values are not wanted (Tomas Vondra)
This is accomplished by setting the new TIMING option to FALSE.

**E.109.3.6. Data Types**

- Add support for range data types (Jeff Davis, Tom Lane, Alexander Korotkov)
  A range data type stores a lower and upper bound belonging to its base data type. It supports operations like contains, overlaps, and intersection.

- Add a JSON data type (Robert Haas)
  This type stores JSON (JavaScript Object Notation) data with proper validation.

- Add `array_to_json()` and `row_to_json()` (Andrew Dunstan)

- Add a SMALLSERIAL data type (Mike Pultz)
  This is like SERIAL, except it stores the sequence in a two-byte integer column (int2).

- Allow domains to be declared NOT VALID (Álvaro Herrera)
  This option can be set at domain creation time, or via `ALTER DOMAIN ... ADD CONSTRAINT ... NOT VALID`.
  `ALTER DOMAIN ... VALIDATE CONSTRAINT` fully validates the constraint.

- Support more locale-specific formatting options for the money data type (Tom Lane)
  Specifically, honor all the POSIX options for ordering of the value, sign, and currency symbol in monetary output. Also, make sure that the thousands separator is only inserted to the left of the decimal point, as required by POSIX.

- Add bitwise “and”, “or”, and “not” operators for the macaddr data type (Brendan Jurd)

- Allow `xpath()` to return a single-element XML array when supplied a scalar value (Florian Pflug)
  Previously, it returned an empty array. This change will also cause `xpath_exists()` to return true, not false, for such expressions.

- Improve XML error handling to be more robust (Florian Pflug)

**E.109.3.7. Functions**

- Allow non-superusers to use `pg_cancel_backend()` and `pg_terminate_backend()` on other sessions belonging to the same user (Magnus Hagander, Josh Kupershmidt, Dan Farina)
  Previously only superusers were allowed to use these functions.

- Allow importing and exporting of transaction snapshots (Joachim Wieland, Tom Lane)
  This allows multiple transactions to share identical views of the database state. Snapshots are exported via `pg_export_snapshot()` and imported via `SET TRANSACTION SNAPSHOT`.
  Only snapshots from currently-running transactions can be imported.

- Support COLLATION FOR on expressions (Peter Eisentraut)
  This returns a string representing the collation of the expression.

- Add `pg_opfamily_is_visible()` (Josh Kupershmidt)

- Add a numeric variant of `pg_size_pretty()` for use with `pg_xlog_location_diff()` (Fujii Masao)
• Add a `pg_trigger_depth()` function (Kevin Grittner)
  This reports the current trigger call depth.
• Allow `string_agg()` to process `bytea` values (Pavel Stehule)
• Fix regular expressions in which a back-reference occurs within a larger quantified subexpression (Tom Lane)
  For example, `^\(\w+\) ( \1)+$`. Previous releases did not check that the back-reference actually matched the first occurrence.

### E.109.3.8. Information Schema

• Add information schema views `role_udt_grants`, `udt_privileges`, and `user_defined_types` (Peter Eisentraut)
• Add composite-type attributes to the information schema `element_types` view (Peter Eisentraut)
• Implement `interval_type` columns in the information schema (Peter Eisentraut)
  Formerly these columns read as nulls.
• Implement collation-related columns in the information schema `attributes`, `columns`, `domains`, and `element_types` views (Peter Eisentraut)
• Implement the `with_hierarchy` column in the information schema `table_privileges` view (Peter Eisentraut)
• Add display of sequence `USAGE` privileges to information schema (Peter Eisentraut)
• Make the information schema show default privileges (Peter Eisentraut)
  Previously, non-empty default permissions were not represented in the views.

### E.109.3.9. Server-Side Languages

#### E.109.3.9.1. PL/pgSQL Server-Side Language

• Allow the PL/pgSQL `OPEN` cursor command to supply parameters by name (Yeb Havinga)
• Add a `GET STACKED DIAGNOSTICS` PL/pgSQL command to retrieve exception info (Pavel Stehule)
• Speed up PL/pgSQL array assignment by caching type information (Pavel Stehule)
• Improve performance and memory consumption for long chains of `ELSIF` clauses (Tom Lane)
• Output the function signature, not just the name, in PL/pgSQL error messages (Pavel Stehule)

#### E.109.3.9.2. PL/Python Server-Side Language

• Add PL/Python SPI cursor support (Jan Urbanski)
  This allows PL/Python to read partial result sets.
• Add result metadata functions to PL/Python (Peter Eisentraut)
  Specifically, this adds result object functions `.colnames`, `.coltypes`, and `.coltypmods.`
• Remove support for Python 2.2 (Peter Eisentraut)

E.109.3.9.3. SQL Server-Side Language

• Allow SQL-language functions to reference parameters by name (Matthew Draper)

To use this, simply name the function arguments and then reference the argument names in the SQL function body.

E.109.3.10. Client Applications

• Add initdb options --auth-local and --auth-host (Peter Eisentraut)

This allows separate control of local and host pg_hba.conf authentication settings. --auth still controls both.

• Add --replication/--no-replication flags to createuser to control replication permission (Fujii Masao)

• Add the --if-exists option to dropdb and dropuser (Josh Kupershmidt)

• Give command-line tools the ability to specify the name of the database to connect to, and fall back to template1 if a postgres database connection fails (Robert Haas)

E.109.3.10.1. psql

• Add a display mode to auto-expand output based on the display width (Peter Eisentraut)

This adds the auto option to the \x command, which switches to the expanded mode when the normal output would be wider than the screen.

• Allow inclusion of a script file that is named relative to the directory of the file from which it was invoked (Gurjeet Singh)

This is done with a new command \ir.

• Add support for non-ASCII characters in psql variable names (Tom Lane)

• Add support for major-version-specific .psqlrc files (Bruce Momjian)

psql already supported minor-version-specific .psqlrc files.

• Provide environment variable overrides for psql history and startup file locations (Andrew Dunstan)

PSQL_HISTORY and PSQLRC now determine these file names if set.

• Add a \setenv command to modify the environment variables passed to child processes (Andrew Dunstan)

• Name psql's temporary editor files with a .sql extension (Peter Eisentraut)

This allows extension-sensitive editors to select the right mode.

• Allow psql to use zero-byte field and record separators (Peter Eisentraut)

Various shell tools use zero-byte (NUL) separators, e.g. find.

• Make the \timing option report times for failed queries (Magnus Hagander)

Previously times were reported only for successful queries.

• Unify and tighten psql's treatment of \copy and SQL COPY (Noah Misch)
This fix makes failure behavior more predictable and honors `\set ON_ERROR_ROLLBACK`.

E.109.3.10.2. Informational Commands

- Make `\d` on a sequence show the table/column name owning it (Magnus Hagander)
- Show statistics target for columns in `\d+` (Magnus Hagander)
- Show role password expiration dates in `\du` (Fabrízio de Royes Mello)
- Display comments for casts, conversions, domains, and languages (Josh Kupershmidt)
  These are included in the output of `\dc+`, `\dc+`, `\dd+`, and `\dl` respectively.
- Display comments for SQL/MED objects (Josh Kupershmidt)
  These are included in the output of `\des+`, `\det+`, and `\dew+` for foreign servers, foreign tables, and foreign data wrappers respectively.
- Change `\dd` to display comments only for object types without their own backslash command (Josh Kupershmidt)

E.109.3.10.3. Tab Completion

- In psql tab completion, complete SQL keywords in either upper or lower case according to the new `COMP_KEYWORD_CASE` setting (Peter Eisentraut)
- Add tab completion support for `EXECUTE` (Andreas Karlsson)
- Allow tab completion of role references in `GRANT/REVOKE` (Peter Eisentraut)
- Allow tab completion of file names to supply quotes, when necessary (Noah Misch)
- Change tab completion support for `TABLE` to also include views (Magnus Hagander)

E.109.3.10.4. pg_dump

- Add an `--exclude-table-data` option to `pg_dump` (Andrew Dunstan)
  This allows dumping of a table’s definition but not its data, on a per-table basis.
- Add a `--section` option to `pg_dump` and `pg_restore` (Andrew Dunstan)
  Valid values are `pre-data`, `data`, and `post-data`. The option can be given more than once to select two or more sections.
- Make `pg_dumpall` dump all roles first, then all configuration settings on roles (Phil Sorber)
  This allows a role’s configuration settings to mention other roles without generating an error.
- Allow `pg_dumpall` to avoid errors if the `postgres` database is missing in the new cluster (Robert Haas)
- Dump foreign server user mappings in user name order (Peter Eisentraut)
  This helps produce deterministic dump files.
- Dump operators in a predictable order (Peter Eisentraut)
- Tighten rules for when extension configuration tables are dumped by `pg_dump` (Tom Lane)
- Make `pg_dump` emit more useful dependency information (Tom Lane)
The dependency links included in archive-format dumps were formerly of very limited use, because they frequently referenced objects that appeared nowhere in the dump. Now they represent actual dependencies (possibly indirect) among the dumped objects.

- Improve pg_dump's performance when dumping many database objects (Tom Lane)

### E.109.3.11. libpq

- Allow libpq connection strings to have the format of a URI (Alexander Shulgin)

  The syntax begins with `postgres://`. This can allow applications to avoid implementing their own parser for URIs representing database connections.

- Add a connection option to disable SSL compression (Laurenz Albe)

  This can be used to remove the overhead of SSL compression on fast networks.

- Add a single-row processing mode for better handling of large result sets (Kyotaro Horiguchi, Marko Kreen)

  Previously, libpq always collected the entire query result in memory before passing it back to the application.

- Add `const` qualifiers to the declarations of the functions `PQconnectdbParams`, `PQconnectStartParams`, and `PQpingParams` (Lionel Elie Mamane)

- Allow the `.pgpass` file to include escaped characters in the password field (Robert Haas)

- Make library functions use `abort()` instead of `exit()` when it is necessary to terminate the process (Peter Eisentraut)

  This choice does not interfere with the normal exit codes used by the program, and generates a signal that can be caught by the caller.

### E.109.3.12. Source Code

- Remove dead ports (Peter Eisentraut)

  The following platforms are no longer supported: dgux, nextstep, sunos4, svr4, ultras4, univel, bsdi.

- Add support for building with MS Visual Studio 2010 (Brar Piening)

- Enable compiling with the MinGW-w64 32-bit compiler (Lars Kanis)

- Install `plpgsql.h` into `include/server` during installation (Heikki Linnakangas)

- Improve the latch facility to include detection of postmaster death (Peter Geoghegan, Heikki Linnakangas, Tom Lane)

  This eliminates one of the main reasons that background processes formerly had to wake up to poll for events.

- Use C flexible array members, where supported (Peter Eisentraut)

- Improve the concurrent transaction regression tests (isolationtester) (Noah Misch)

- Modify thread_test to create its test files in the current directory, rather than `/tmp` (Bruce Momjian)

- Improve flex and bison warning and error reporting (Tom Lane)

- Add memory barrier support (Robert Haas)
This is currently unused.

• Modify pgindent to use a typedef file (Bruce Momjian)
• Add a hook for processing messages due to be sent to the server log (Martin Pihlak)
• Add object access hooks for DROP commands (KaiGai Kohei)
• Centralize DROP handling for some object types (KaiGai Kohei)
• Add a pg_upgrade test suite (Peter Eisentraut)
• Sync regular expression code with TCL 8.5.11 and improve internal processing (Tom Lane)
• Move CRC tables to libpqport, and provide them in a separate include file (Daniel Farina)
• Add options to git_changelog for use in major release note creation (Bruce Momjian)
• Support Linux's /proc/self/oom_score_adj API (Tom Lane)

E.109.3.13. Additional Modules

• Improve efficiency of dblink by using libpq's new single-row processing mode (Kyotaro Horiguchi, Marko Kreen)

This improvement does not apply to dblink_send_query()/dblink_get_result().

• Support force_not_null option in file_fdw (Shigeru Hanada)
• Implement dry-run mode for pg_archivecleanup (Gabriele Bartolini)

This only outputs the names of files to be deleted.

• Add new pgbench switches --unlogged-tables, --tablespace, and --index-tablespace (Robert Haas)
• Change pg_test_fsync to test for a fixed amount of time, rather than a fixed number of cycles (Bruce Momjian)

The -o/cycles option was removed, and -s/seconds added.

• Add a pg_test_timing utility to measure clock monotonicity and timing overhead (Ants Aasma, Greg Smith)
• Add a tcn (triggered change notification) module to generate NOTIFY events on table changes (Kevin Grittner)

E.109.3.13.1. pg_upgrade

• Adjust pg_upgrade environment variables (Bruce Momjian)

Rename data, bin, and port environment variables to begin with PG, and support PGPORTOLD/PGPORTNEW, to replace PGPORT.

• Overhaul pg_upgrade logging and failure reporting (Bruce Momjian)

Create four append-only log files, and delete them on success. Add -r/--retain option to unconditionally retain these files. Also remove pg_upgrade options -g/-G/-l options as unnecessary, and tighten log file permissions.
• Make pg_upgrade create a script to incrementally generate more accurate optimizer statistics (Bruce Momjian)
This reduces the time needed to generate minimal cluster statistics after an upgrade.

- Allow `pg_upgrade` to upgrade an old cluster that does not have a `postgres` database (Bruce Momjian)
- Allow `pg_upgrade` to handle cases where some old or new databases are missing, as long as they are empty (Bruce Momjian)
- Allow `pg_upgrade` to handle configuration-only directory installations (Bruce Momjian)
- In `pg_upgrade`, add `-o/-O` options to pass parameters to the servers (Bruce Momjian)
  This is useful for configuration-only directory installs.
- Change `pg_upgrade` to use port 50432 by default (Bruce Momjian)
  This helps avoid unintended client connections during the upgrade.
- Reduce cluster locking in `pg_upgrade` (Bruce Momjian)
  Specifically, only lock the old cluster if link mode is used, and do it right after the schema is restored.

**E.109.3.13.2. pg_stat_statements**

- Allow `pg_stat_statements` to aggregate similar queries via SQL text normalization (Peter Geoghegan, Tom Lane)
  Users with applications that use non-parameterized SQL will now be able to monitor query performance without detailed log analysis.
- Add dirtied and written block counts and read/write times to `pg_stat_statements` (Robert Haas, Ants Aasma)
- Prevent `pg_stat_statements` from double-counting `PREPARE` and `EXECUTE` commands (Tom Lane)

**E.109.3.13.3. sepgsql**

- Support `SECURITY LABEL` on global objects (KaiGai Kohei, Robert Haas)
  Specifically, add security labels to databases, tablespaces, and roles.
- Allow sepgsql to honor database labels (KaiGai Kohei)
- Perform sepgsql permission checks during the creation of various objects (KaiGai Kohei)
- Add `sepgsql_setcon()` and related functions to control the sepgsql security domain (KaiGai Kohei)
- Add a user space access cache to sepgsql to improve performance (KaiGai Kohei)

**E.109.3.14. Documentation**

- Add a rule to optionally build HTML documentation using the stylesheet from the website (Magnus Hagander)
  Use `gmake STYLE=website draft`.
- Improve `EXPLAIN` documentation (Tom Lane)
- Document that user/database names are preserved with double-quot ing by command-line tools like `vacuumdb` (Bruce Momjian)
• Document the actual string returned by the client for MD5 authentication (Cyan Ogilvie)

• Deprecate use of GLOBAL and LOCAL in CREATE TEMP TABLE (Noah Misch)

PostgreSQL has long treated these keywords as no-ops, and continues to do so; but in future they might mean what the SQL standard says they mean, so applications should avoid using them.

E.110. Release 9.1.24

Release date: 2016-10-27

This release contains a variety of fixes from 9.1.23. For information about new features in the 9.1 major release, see Section E.134.

This is expected to be the last PostgreSQL release in the 9.1.X series. Users are encouraged to update to a newer release branch soon.

E.110.1. Migration to Version 9.1.24

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.110.2. Changes

• Fix EvalPlanQual rechecks involving CTE scans (Tom Lane)

The recheck would always see the CTE as returning no rows, typically leading to failure to update rows that were recently updated.

• Fix improper repetition of previous results from hashed aggregation in a subquery (Andrew Gierth)

The test to see if we can reuse a previously-computed hash table of the aggregate state values neglected the possibility of an outer query reference appearing in an aggregate argument expression. A change in the value of such a reference should lead to recalculating the hash table, but did not.

• Fix timeout length when VACUUM is waiting for exclusive table lock so that it can truncate the table (Simon Riggs)

The timeout was meant to be 50 milliseconds, but it was actually only 50 microseconds, causing VACUUM to give up on truncation much more easily than intended. Set it to the intended value.

• Remove artificial restrictions on the values accepted by numeric_in() and numeric_recv() (Tom Lane)

We allow numeric values up to the limit of the storage format (more than 1e100000), so it seems fairly pointless that numeric_in() rejected scientific-notation exponents above 1000. Likewise, it was silly for numeric_recv() to reject more than 1000 digits in an input value.

• Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)

• Fix file descriptor leakage when truncating a temporary relation of more than 1GB (Andres Freund)

•Disallow starting a standalone backend with standby_mode turned on (Michael Paquier)

This can't do anything useful, since there will be no WAL receiver process to fetch more WAL data; and it could result in misbehavior in code that wasn't designed with this situation in mind.
• Don’t try to share SSL contexts across multiple connections in libpq (Heikki Linnakangas)

This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

• Avoid corner-case memory leak in libpq (Tom Lane)

The reported problem involved leaking an error report during PQreset(), but there might be related cases.

• Make ecpg's --help and --version options work consistently with our other executables (Haribabu Kommi)

• Fix contrib/intarray/bench/bench.pl to print the results of the EXPLAIN it does when given the -e option (Daniel Gustafsson)

• Prevent failure of obsolete dynamic time zone abbreviations (Tom Lane)

If a dynamic time zone abbreviation does not match any entry in the referenced time zone, treat it as equivalent to the time zone name. This avoids unexpected failures when IANA removes abbreviations from their time zone database, as they did in tzdata release 2016f and seem likely to do again in the future. The consequences were not limited to not recognizing the individual abbreviation; any mismatch caused the pg_timezone_abbrevs view to fail altogether.

• Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.111. Release 9.1.23

Release date: 2016-08-11

This release contains a variety of fixes from 9.1.22. For information about new features in the 9.1 major release, see Section E.134.

The PostgreSQL community will stop releasing updates for the 9.1.X release series in September 2016. Users are encouraged to update to a newer release branch soon.

E.111.1. Migration to Version 9.1.23

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.111.2. Changes

• Fix possible mis-evaluation of nested CASE-WHEN expressions (Heikki Linnakangas, Michael Paquier, Tom Lane)
A CASE expression appearing within the test value subexpression of another CASE could become confused about whether its own test value was null or not. Also, inlining of a SQL function implementing the equality operator used by a CASE expression could result in passing the wrong test value to functions called within a CASE expression in the SQL function's body. If the test values were of different data types, a crash might result; moreover such situations could be abused to allow disclosure of portions of server memory. (CVE-2016-5423)

- Fix client programs' handling of special characters in database and role names (Noah Misch, Nathan Bossart, Michael Paquier)

Numerous places in vacuumdb and other client programs could become confused by database and role names containing double quotes or backslashes. Tighten up quoting rules to make that safe. Also, ensure that when a conninfo string is used as a database name parameter to these programs, it is correctly treated as such throughout.

Fix handling of paired double quotes in psql's `\connect` and `\password` commands to match the documentation.

Introduce a new `-reuse-previous` option in psql's `\connect` command to allow explicit control of whether to re-use connection parameters from a previous connection. (Without this, the choice is based on whether the database name looks like a conninfo string, as before.) This allows secure handling of database names containing special characters in pg_dumpall scripts.

pg_dumpall now refuses to deal with database and role names containing carriage returns or newlines, as it seems impractical to quote those characters safely on Windows. In future we may reject such names on the server side, but that step has not been taken yet.

These are considered security fixes because crafted object names containing special characters could have been used to execute commands with superuser privileges the next time a superuser executes pg_dumpall or other routine maintenance operations. (CVE-2016-5424)

- Fix corner-case misbehaviors for `IS NULL/IS NOT NULL` applied to nested composite values (Andrew Gierth, Tom Lane)

The SQL standard specifies that `IS NULL` should return TRUE for a row of all null values (thus `ROW(NULL,NULL) IS NULL` yields TRUE), but this is not meant to apply recursively (thus `ROW(NULL, ROW(NULL,NULL)) IS NULL` yields FALSE). The core executor got this right, but certain planner optimizations treated the test as recursive (thus producing TRUE in both cases), and contrib/postgres_fdw could produce remote queries that misbehaved similarly.

- Make the `inet` and `cidr` data types properly reject IPv6 addresses with too many colon-separated fields (Tom Lane)

- Prevent crash in `close_ps()` (the point `## lseg` operator) for NaN input coordinates (Tom Lane)

Make it return NULL instead of crashing.

- Fix several one-byte buffer over-reads in `to_number()` (Peter Eisentraut)

In several cases the `to_number()` function would read one more character than it should from the input string. There is a small chance of a crash, if the input happens to be adjacent to the end of memory.

- Avoid unsafe intermediate state during expensive paths through `heap_update()` (Masahiko Sawada, Andres Freund)

Previously, these cases locked the target tuple (by setting its XMAX) but did not WAL-log that action, thus risking data integrity problems if the page were spilled to disk and then a database crash occurred before the tuple update could be completed.
• Avoid consuming a transaction ID during VACUUM (Alexander Korotkov)

Some cases in VACUUM unnecessarily caused an XID to be assigned to the current transaction. Normally this is negligible, but if one is up against the XID wraparound limit, consuming more XIDs during anti-wraparound vacuums is a very bad thing.

• Avoid canceling hot-standby queries during VACUUM FREEZE (Simon Riggs, Álvaro Herrera)

VACUUM FREEZE on an otherwise-idle master server could result in unnecessary cancellations of queries on its standby servers.

• When a manual ANALYZE specifies a column list, don't reset the table's changes_since_analyze counter (Tom Lane)

If we're only analyzing some columns, we should not prevent routine auto-analyze from happening for the other columns.

• Fix ANALYZE's overestimation of n_distinct for a unique or nearly-unique column with many null entries (Tom Lane)

The nulls could get counted as though they were themselves distinct values, leading to serious planner misestimates in some types of queries.

• Prevent autovacuum from starting multiple workers for the same shared catalog (Álvaro Herrera)

Normally this isn't much of a problem because the vacuum doesn't take long anyway; but in the case of a severely bloated catalog, it could result in all but one worker uselessly waiting instead of doing useful work on other tables.

• Fix contrib/btree_gin to handle the smallest possible bigint value correctly (Peter Eisentraut)

• Teach libpq to correctly decode server version from future servers (Peter Eisentraut)

It's planned to switch to two-part instead of three-part server version numbers for releases after 9.6. Make sure that PQserverVersion() returns the correct value for such cases.

• Fix ecpg's code for unsigned long long array elements (Michael Meskes)

• Make pg_basebackup accept -Z 0 as specifying no compression (Fujii Masao)

• Revert to the old heuristic timeout for pg_ctl start -w (Tom Lane)

The new method adopted as of release 9.1.20 does not work when silent_mode is enabled, so go back to the old way.

• Fix makefiles' rule for building AIX shared libraries to be safe for parallel make (Noah Misch)

• Fix TAP tests and MSVC scripts to work when build directory's path name contains spaces (Michael Paquier, Kyotaro Horiguchi)

• Make regression tests safe for Danish and Welsh locales (Jeff Janes, Tom Lane)

Change some test data that triggered the unusual sorting rules of these locales.

• Update our copy of the timezone code to match IANA's tzcode release 2016c (Tom Lane)

This is needed to cope with anticipated future changes in the time zone data files. It also fixes some corner-case bugs in coping with unusual time zones.

• Update time zone data files to tzdata release 2016f for DST law changes in Kemerovo and Novosibirsk, plus historical corrections for Azerbaijan, Belarus, and Morocco.
E.112. Release 9.1.22

Release date: 2016-05-12

This release contains a variety of fixes from 9.1.21. For information about new features in the 9.1 major release, see Section E.134.

The PostgreSQL community will stop releasing updates for the 9.1.X release series in September 2016. Users are encouraged to update to a newer release branch soon.

E.112.1. Migration to Version 9.1.22

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.112.2. Changes

- Clear the OpenSSL error queue before OpenSSL calls, rather than assuming it's clear already; and make sure we leave it clear afterwards (Peter Geoghegan, Dave Vitke, Peter Eisentraut)

  This change prevents problems when there are multiple connections using OpenSSL within a single process and not all the code involved follows the same rules for when to clear the error queue. Failures have been reported specifically when a client application uses SSL connections in libpq concurrently with SSL connections using the PHP, Python, or Ruby wrappers for OpenSSL. It's possible for similar problems to arise within the server as well, if an extension module establishes an outgoing SSL connection.

- Fix "failed to build any N-way joins" planner error with a full join enclosed in the right-hand side of a left join (Tom Lane)

- Fix possible misbehavior of TH, th, and Y, YYY format codes in to_timestamp() (Tom Lane)

  These could advance off the end of the input string, causing subsequent format codes to read garbage.

- Fix dumping of rules and views in which the array argument of a value operator ANY (array) construct is a sub-SELECT (Tom Lane)

- Make pg_regress use a startup timeout from the PGCTLTIMEOUT environment variable, if that's set (Tom Lane)

  This is for consistency with a behavior recently added to pg_ctl; it eases automated testing on slow machines.

- Fix pg_upgrade to correctly restore extension membership for operator families containing only one operator class (Tom Lane)

  In such a case, the operator family was restored into the new database, but it was no longer marked as part of the extension. This had no immediate ill effects, but would cause later pg_dump runs to emit output that would cause (harmless) errors on restore.

- Rename internal function strtol() to strtoint() to avoid conflict with a NetBSD library function (Thomas Munro)

- Fix reporting of errors from bind() and listen() system calls on Windows (Tom Lane)

- Reduce verbosity of compiler output when building with Microsoft Visual Studio (Christian Ullrich)

- Avoid possibly-unsafe use of Windows’ FormatMessage() function (Christian Ullrich)
Use the FORMAT_MESSAGE_IGNORE_INSERTS flag where appropriate. No live bug is known to exist here, but it seems like a good idea to be careful.

- Update time zone data files to tzdata release 2016d for DST law changes in Russia and Venezuela. There are new zone names Europe/Kirov and Asia/Tomsk to reflect the fact that these regions now have different time zone histories from adjacent regions.

E.113. Release 9.1.21

Release date: 2016-03-31

This release contains a variety of fixes from 9.1.20. For information about new features in the 9.1 major release, see Section E.134.

E.113.1. Migration to Version 9.1.21

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.113.2. Changes

- Fix incorrect handling of NULL index entries in indexed \texttt{ROW()} comparisons (Tom Lane)

  An index search using a row comparison such as \texttt{ROW(a, b) > ROW('x', 'y')} would stop upon reaching a NULL entry in the \texttt{b} column, ignoring the fact that there might be non-NULL \texttt{b} values associated with later values of \texttt{a}.

- Avoid unlikely data-loss scenarios due to renaming files without adequate \texttt{fsync()} calls before and after (Michael Paquier, Tomas Vondra, Andres Freund)

- Correctly handle cases where \texttt{pg_subtrans} is close to XID wraparound during server startup (Jeff Janes)

- Fix corner-case crash due to trying to free \texttt{localeconv()} output strings more than once (Tom Lane)

- Fix parsing of affix files for \texttt{ispell} dictionaries (Tom Lane)

  The code could go wrong if the affix file contained any characters whose byte length changes during case-folding, for example \texttt{I} in Turkish UTF8 locales.

- Avoid use of \texttt{sscanf()} to parse \texttt{ispell} dictionary files (Artur Zakirov)

  This dodges a portability problem on FreeBSD-derived platforms (including macOS).

- Avoid a crash on old Windows versions (before 7SP1/2008R2SP1) with an AVX2-capable CPU and a Postgres build done with Visual Studio 2013 (Christian Ullrich)

  This is a workaround for a bug in Visual Studio 2013's runtime library, which Microsoft have stated they will not fix in that version.

- Fix psql's tab completion logic to handle multibyte characters properly (Kyotaro Horiguchi, Robert Haas)

- Fix psql's tab completion for \texttt{SECURITY LABEL} (Tom Lane)

  Pressing TAB after \texttt{SECURITY LABEL} might cause a crash or offering of inappropriate keywords.
• Make pg_ctl accept a wait timeout from the PGCTLTIMEOUT environment variable, if none is specified on the command line (Noah Misch)

This eases testing of slower buildfarm members by allowing them to globally specify a longer-than-normal timeout for postmaster startup and shutdown.

• Fix incorrect test for Windows service status in pg_ctl (Manuel Mathar)

The previous set of minor releases attempted to fix pg_ctl to properly determine whether to send log messages to Window’s Event Log, but got the test backwards.

• Fix pgbench to correctly handle the combination of -c and -M prepared options (Tom Lane)

• In PL/Perl, properly translate empty Postgres arrays into empty Perl arrays (Alex Hunsaker)

• Make PL/Python cope with function names that aren’t valid Python identifiers (Jim Nasby)

• Fix multiple mistakes in the statistics returned by contrib/pgstatupid's pgstatindex() function (Tom Lane)

• Remove dependency on psemd in MSVC builds, since it’s no longer provided by core Perl (Michael Paquier, Andrew Dunstan)

• Update time zone data files to tzdata release 2016c for DST law changes in Azerbaijan, Chile, Haiti, Palestine, and Russia (Altai, Astrakhan, Kirov, Sakhalin, Ulyanovsk regions), plus historical corrections for Lithuania, Moldova, and Russia (Kaliningrad, Samara, Volgograd).

E.114. Release 9.1.20

Release date: 2016-02-11

This release contains a variety of fixes from 9.1.19. For information about new features in the 9.1 major release, see Section E.134.

E.114.1. Migration to Version 9.1.20

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.114.2. Changes

• Fix infinite loops and buffer-overrun problems in regular expressions (Tom Lane)

Very large character ranges in bracket expressions could cause infinite loops in some cases, and memory overwrites in other cases. (CVE-2016-0773)

• Perform an immediate shutdown if the postmaster.pid file is removed (Tom Lane)

The postmaster now checks every minute or so that postmaster.pid is still there and still contains its own PID. If not, it performs an immediate shutdown, as though it had received SIGQUIT. The main motivation for this change is to ensure that failed buildfarm runs will get cleaned up without manual intervention; but it also serves to limit the bad effects if a DBA forcibly removes postmaster.pid and then starts a new postmaster.

• In SERIALIZABLE transaction isolation mode, serialization anomalies could be missed due to race conditions during insertions (Kevin Grittner, Thomas Munro)

• Fix failure to emit appropriate WAL records when doing ALTER TABLE ... SET TABLESPACE for unlogged relations (Michael Paquier, Andres Freund)
Even though the relation's data is unlogged, the move must be logged or the relation will be inaccessible after a standby is promoted to master.

- Fix possible misinitialization of unlogged relations at the end of crash recovery (Andres Freund, Michael Paquier)
- Fix `ALTER COLUMN TYPE` to reconstruct inherited check constraints properly (Tom Lane)
- Fix `REASSIGN OWNED` to change ownership of composite types properly (Álvaro Herrera)
- Fix `REASSIGN OWNED` and `ALTER OWNER` to correctly update granted-permissions lists when changing owners of data types, foreign data wrappers, or foreign servers (Bruce Momjian, Álvaro Herrera)
- Fix `REASSIGN OWNED` to ignore foreign user mappings, rather than fail (Álvaro Herrera)
- Add more defenses against bad planner cost estimates for GIN index scans when the index's internal statistics are very out-of-date (Tom Lane)
- Make planner cope with hypothetical GIN indexes suggested by an index advisor plug-in (Julien Rouhaud)
- Fix dumping of whole-row Vars in `ROW()` and `VALUES()` lists (Tom Lane)
- Fix possible internal overflow in numeric division (Dean Rasheed)
- Fix enforcement of restrictions inside parentheses within regular expression lookahead constraints (Tom Lane)

Lookahead constraints aren't allowed to contain backrefs, and parentheses within them are always considered non-capturing, according to the manual. However, the code failed to handle these cases properly inside a parenthesized subexpression, and would give unexpected results.

- Conversion of regular expressions to indexscan bounds could produce incorrect bounds from regexps containing lookahead constraints (Tom Lane)
- Fix regular-expression compiler to handle loops of constraint arcs (Tom Lane)

The code added for CVE-2007-4772 was both incomplete, in that it didn't handle loops involving more than one state, and incorrect, in that it could cause assertion failures (though there seem to be no bad consequences of that in a non-assert build). Multi-state loops would cause the compiler to run until the query was canceled or it reached the too-many-states error condition.

- Improve memory-usage accounting in regular-expression compiler (Tom Lane)

This causes the code to emit “regular expression is too complex” errors in some cases that previously used unreasonable amounts of time and memory.

- Improve performance of regular-expression compiler (Tom Lane)
- Make `%h` and `%r` escapes in `log_line_prefix` work for messages emitted due to `log_connections` (Tom Lane)

Previously, `%h/%r` started to work just after a new session had emitted the “connection received” log message; now they work for that message too.

- On Windows, ensure the shared-memory mapping handle gets closed in child processes that don’t need it (Tom Lane, Amit Kapila)

This oversight resulted in failure to recover from crashes whenever `logging_collector` is turned on.
• Fix possible failure to detect socket EOF in non-blocking mode on Windows (Tom Lane)

It's not entirely clear whether this problem can happen in pre-9.5 branches, but if it did, the symptom
would be that a walsender process would wait indefinitely rather than noticing a loss of connection.

• Avoid leaking a token handle during SSPI authentication (Christian Ullrich)

• In psql, ensure that libreadline's idea of the screen size is updated when the terminal window size
changes (Merlin Moncure)

Previously, libreadline did not notice if the window was resized during query output, leading to
strange behavior during later input of multiline queries.

• Fix psql's \det command to interpret its pattern argument the same way as other \d commands
with potentially schema-qualified patterns do (Reece Hart)

• Avoid possible crash in psql's \c command when previous connection was via Unix socket and
command specifies a new hostname and same username (Tom Lane)

• In pg_ctl start -w, test child process status directly rather than relying on heuristics (Tom
Lane, Michael Paquier)

Previously, pg_ctl relied on an assumption that the new postmaster would always create post-
master.pid within five seconds. But that can fail on heavily-loaded systems, causing pg_ctl to
report incorrectly that the postmaster failed to start.

Except on Windows, this change also means that a pg_ctl start -w done immediately after
another such command will now reliably fail, whereas previously it would report success if done
within two seconds of the first command.

• In pg_ctl start -w, don't attempt to use a wildcard listen address to connect to the postmaster
(Kondo Yuta)

On Windows, pg_ctl would fail to detect postmaster startup if listen_addresses is set to
0.0.0.0 or ::, because it would try to use that value verbatim as the address to connect to, which
doesn't work. Instead assume that 127.0.0.1 or ::1, respectively, is the right thing to use.

• In pg_ctl on Windows, check service status to decide where to send output, rather than checking if
standard output is a terminal (Michael Paquier)

• In pg_dump and pg_basebackup, adopt the GNU convention for handling tar-archive members ex-
ceeding 8GB (Tom Lane)

The POSIX standard for tar file format does not allow archive member files to exceed 8GB, but
most modern implementations of tar support an extension that fixes that. Adopt this extension so that
pg_dump with -Ft no longer fails on tables with more than 8GB of data, and so that pg_basebackup
can handle files larger than 8GB. In addition, fix some portability issues that could cause failures
for members between 4GB and 8GB on some platforms. Potentially these problems could cause
unrecoverable data loss due to unreadable backup files.

• Fix assorted corner-case bugs in pg_dump's processing of extension member objects (Tom Lane)

• Make pg_dump mark a view's triggers as needing to be processed after its rule, to prevent possible
failure during parallel pg_restore (Tom Lane)

• Ensure that relation option values are properly quoted in pg_dump (Kouhei Sutou, Tom Lane)

A reloption value that isn't a simple identifier or number could lead to dump/reload failures due
to syntax errors in CREATE statements issued by pg_dump. This is not an issue with any relo-
option currently supported by core PostgreSQL, but extensions could allow relocations that cause the
problem.
Fix pg_upgrade's file-copying code to handle errors properly on Windows (Bruce Momjian)

Install guards in pgbench against corner-case overflow conditions during evaluation of script-specified division or modulo operators (Fabien Coelho, Michael Paquier)

Prevent certain PL/Java parameters from being set by non-superusers (Noah Misch)

This change mitigates a PL/Java security bug (CVE-2016-0766), which was fixed in PL/Java by marking these parameters as superuser-only. To fix the security hazard for sites that update PostgreSQL more frequently than PL/Java, make the core code aware of them also.

Improve libpq's handling of out-of-memory situations (Michael Paquier, Amit Kapila, Heikki Linnakangas)

Fix order of arguments in ecpg-generated typedef statements (Michael Meskes)

Use %g not %f format in ecpg's PGTYPESnumeric_from_double() (Tom Lane)

Fix ecpg-supplied header files to not contain comments continued from a preprocessor directive line onto the next line (Michael Meskes)

Such a comment is rejected by ecpg. It's not yet clear whether ecpg itself should be changed.

Ensure that contrib/pgcrypto's crypt() function can be interrupted by query cancel (Andreas Karlsson)

Accept flex versions later than 2.5.x (Tom Lane, Michael Paquier)

Now that flex 2.6.0 has been released, the version checks in our build scripts needed to be adjusted.

Install our missing script where PGXS builds can find it (Jim Nasby)

This allows sane behavior in a PGXS build done on a machine where build tools such as bison are missing.

Ensure that dynloader.h is included in the installed header files in MSVC builds (Bruce Momjian, Michael Paquier)

Add variant regression test expected-output file to match behavior of current libxml2 (Tom Lane)

The fix for libxml2's CVE-2015-7499 causes it not to output error context reports in some cases where it used to do so. This seems to be a bug, but we'll probably have to live with it for some time, so work around it.

Update time zone data files to tzdata release 2016a for DST law changes in Cayman Islands, Metlakatla, and Trans-Baikal Territory (Zabaykalsky Krai), plus historical corrections for Pakistan.

### E.115. Release 9.1.19

**Release date:** 2015-10-08

This release contains a variety of fixes from 9.1.18. For information about new features in the 9.1 major release, see Section E.134.

#### E.115.1. Migration to Version 9.1.19

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.
### E.115.2. Changes

- **Fix contrib/pgcrypto** to detect and report too-short `crypt()` salts (Josh Kupershmidt)

  Certain invalid salt arguments crashed the server or disclosed a few bytes of server memory. We have not ruled out the viability of attacks that arrange for presence of confidential information in the disclosed bytes, but they seem unlikely. (CVE-2015-5288)

- **Fix subtransaction cleanup after a portal (cursor) belonging to an outer subtransaction fails** (Tom Lane, Michael Paquier)

  A function executed in an outer-subtransaction cursor could cause an assertion failure or crash by referencing a relation created within an inner subtransaction.

- **Fix insertion of relations into the relation cache “init file”** (Tom Lane)

  An oversight in a patch in the most recent minor releases caused `pg_trigger_tgrelid_tg-name_index` to be omitted from the init file. Subsequent sessions detected this, then deemed the init file to be broken and silently ignored it, resulting in a significant degradation in session startup time. In addition to fixing the bug, install some guards so that any similar future mistake will be more obvious.

- **Avoid O(N^2) behavior when inserting many tuples into a SPI query result** (Neil Conway)

- **Improve LISTEN startup time when there are many unread notifications** (Matt Newell)

- **Back-patch 9.3-era addition of per-resource-owner lock caches** (Jeff Janes)

  This substantially improves performance when `pg_dump` tries to dump a large number of tables.

- **Disable SSL renegotiation by default** (Michael Paquier, Andres Freund)

  While use of SSL renegotiation is a good idea in theory, we have seen too many bugs in practice, both in the underlying OpenSSL library and in our usage of it. Renegotiation will be removed entirely in 9.5 and later. In the older branches, just change the default value of `ssl_renegotiation_limit` to zero (disabled).

- **Lower the minimum values of the *_freeze_max_age* parameters** (Andres Freund)

  This is mainly to make tests of related behavior less time-consuming, but it may also be of value for installations with limited disk space.

- **Limit the maximum value of `wal_buffers` to 2GB to avoid server crashes** (Josh Berkus)

- **Fix rare internal overflow in multiplication of numeric values** (Dean Rasheed)

- **Guard against hard-to-reach stack overflows involving record types, range types, json, jsonb, tsquery, ltxtquery and query_int** (Noah Misch)

- **Fix handling of DOW and DOY in datetime input** (Greg Stark)

  These tokens aren’t meant to be used in datetime values, but previously they resulted in opaque internal error messages rather than “invalid input syntax”.

- **Add more query-cancel checks to regular expression matching** (Tom Lane)

- **Add recursion depth protections to regular expression, SIMILAR TO, and LIKE matching** (Tom Lane)

  Suitable search patterns and a low stack depth limit could lead to stack-overflow crashes.

- **Fix potential infinite loop in regular expression execution** (Tom Lane)
A search pattern that can apparently match a zero-length string, but actually doesn't match because of a back reference, could lead to an infinite loop.

- Fix low-memory failures in regular expression compilation (Andreas Seltenreich)
- Fix low-probability memory leak during regular expression execution (Tom Lane)
- Fix rare low-memory failure in lock cleanup during transaction abort (Tom Lane)
- Fix “unexpected out-of-memory situation during sort” errors when using tuplestores with small work_mem settings (Tom Lane)
- Fix very-low-probability stack overrun in qsort (Tom Lane)
- Fix “invalid memory alloc request size” failure in hash joins with large work_mem settings (Tomas Vondra, Tom Lane)
- Fix assorted planner bugs (Tom Lane)

These mistakes could lead to incorrect query plans that would give wrong answers, or to assertion failures in assert-enabled builds, or to odd planner errors such as “could not devise a query plan for the given query”, “could not find pathkey item to sort”, “plan should not reference subplan's variable”, or “failed to assign all NestLoopParams to plan nodes”. Thanks are due to Andreas Seltenreich and Piotr Stefaniak for fuzz testing that exposed these problems.

- Use fuzzy path cost tiebreaking rule in all supported branches (Tom Lane)

This change is meant to avoid platform-specific behavior when alternative plan choices have effectively-identical estimated costs.

- Ensure standby promotion trigger files are removed at postmaster startup (Michael Paquier, Fujii Masao)

This prevents unwanted promotion from occurring if these files appear in a database backup that is used to initialize a new standby server.

- During postmaster shutdown, ensure that per-socket lock files are removed and listen sockets are closed before we remove the postmaster.pid file (Tom Lane)

This avoids race-condition failures if an external script attempts to start a new postmaster as soon as pg_ctl stop returns.

- Fix postmaster's handling of a startup-process crash during crash recovery (Tom Lane)

If, during a crash recovery cycle, the startup process crashes without having restored database consistency, we'd try to launch a new startup process, which typically would just crash again, leading to an infinite loop.

- Do not print a WARNING when an autovacuum worker is already gone when we attempt to signal it, and reduce log verbosity for such signals (Tom Lane)

- Prevent autovacuum launcher from sleeping unduly long if the server clock is moved backwards a large amount (Álvaro Herrera)

- Ensure that cleanup of a GIN index's pending-insertions list is interruptable by cancel requests (Jeff Janes)

- Allow all-zeroes pages in GIN indexes to be reused (Heikki Linnakangas)

Such a page might be left behind after a crash.
• Fix off-by-one error that led to otherwise-harmless warnings about “apparent wraparound” in subtrans/multixact truncation (Thomas Munro)

• Fix misreporting of CONTINUE and MOVE statement types in PL/pgSQL’s error context messages (Pavel Stehule, Tom Lane)

• Fix PL/Perl to handle non-ASCII error message texts correctly (Alex Hunsaker)

• Fix PL/Python crash when returning the string representation of a record result (Tom Lane)

• Fix some places in PL/Tcl that neglected to check for failure of malloc() calls (Michael Paquier, Álvaro Herrera)

• In contrib/isn, fix output of ISBN-13 numbers that begin with 979 (Fabien Coelho)

  EANs beginning with 979 (but not 9790) are considered ISBNs, but they must be printed in the new 13-digit format, not the 10-digit format.

• Improve libpq’s handling of out-of-memory conditions (Michael Paquier, Heikki Linnakangas)

• Fix memory leaks and missing out-of-memory checks in ecpg (Michael Paquier)

• Fix psql’s code for locale-aware formatting of numeric output (Tom Lane)

  The formatting code invoked by \pset numericlocale on did the wrong thing for some uncommon cases such as numbers with an exponent but no decimal point. It could also mangle already-localized output from the money data type.

• Prevent crash in psql’s \c command when there is no current connection (Noah Misch)

• Fix selection of default zlib compression level in pg_dump’s directory output format (Andrew Dunstan)

• Ensure that temporary files created during a pg_dump run with tar-format output are not world-readable (Michael Paquier)

• Fix pg_dump and pg_upgrade to support cases where the postgres or template1 database is in a non-default tablespace (Marti Raudsepp, Bruce Momjian)

• Fix pg_dump to handle object privileges sanely when dumping from a server too old to have a particular privilege type (Tom Lane)

  When dumping functions or procedural languages from pre-7.3 servers, pg_dump would produce GRANT/REVOKE commands that revoked the owner's grantable privileges and instead granted all privileges to PUBLIC. Since the privileges involved are just USAGE and EXECUTE, this isn't a security problem, but it's certainly a surprising representation of the older systems' behavior. Fix it to leave the default privilege state alone in these cases.

• Fix pg_dump to dump shell types (Tom Lane)

  Shell types (that is, not-yet-fully-defined types) aren't useful for much, but nonetheless pg_dump should dump them.

• Fix assorted minor memory leaks in pg_dump and other client-side programs (Michael Paquier)

• Fix spinlock assembly code for PPC hardware to be compatible with AIX’s native assembler (Tom Lane)

  Building with gcc didn't work if gcc had been configured to use the native assembler, which is becoming more common.

• On AIX, test the -qlonglong compiler option rather than just assuming it's safe to use (Noah Misch)
• On AIX, use `-Wl,-brtlib` link option to allow symbols to be resolved at runtime (Noah Misch)

Perl relies on this ability in 5.8.0 and later.

• Avoid use of inline functions when compiling with 32-bit xlc, due to compiler bugs (Noah Misch)

• Use `librt` for `sched_yield()` when necessary, which it is on some Solaris versions (Oskari Saarenmaa)

• Fix Windows `install.bat` script to handle target directory names that contain spaces (Heikki Linnakangas)

• Make the numeric form of the PostgreSQL version number (e.g., `90405`) readily available to extension Makefiles, as a variable named `VERSION_NUM` (Michael Paquier)

• Update time zone data files to `tzdata` release 2015g for DST law changes in Cayman Islands, Fiji, Moldova, Morocco, Norfolk Island, North Korea, Turkey, and Uruguay. There is a new zone name `America/Fort_Nelson` for the Canadian Northern Rockies.

E.116. Release 9.1.18

Release date: 2015-06-12

This release contains a small number of fixes from 9.1.17. For information about new features in the 9.1 major release, see Section E.134.

E.116.1. Migration to Version 9.1.18

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.116.2. Changes

• Fix rare failure to invalidate relation cache init file (Tom Lane)

With just the wrong timing of concurrent activity, a `VACUUM FULL` on a system catalog might fail to update the “init file” that’s used to avoid cache-loading work for new sessions. This would result in later sessions being unable to access that catalog at all. This is a very ancient bug, but it’s so hard to trigger that no reproducible case had been seen until recently.

• Avoid deadlock between incoming sessions and `CREATE/DROP DATABASE` (Tom Lane)

A new session starting in a database that is the target of a `DROP DATABASE` command, or is the template for a `CREATE DATABASE` command, could cause the command to wait for five seconds and then fail, even if the new session would have exited before that.

E.117. Release 9.1.17

Release date: 2015-06-04

This release contains a small number of fixes from 9.1.16. For information about new features in the 9.1 major release, see Section E.134.

E.117.1. Migration to Version 9.1.17

A dump/restore is not required for those running 9.1.X.
However, if you are upgrading from a version earlier than 9.1.16, see Section E.118.

E.117.2. Changes

- Avoid failures while fsync'ing data directory during crash restart (Abhijit Menon-Sen, Tom Lane)

In the previous minor releases we added a patch to fsync everything in the data directory after a crash. Unfortunately its response to any error condition was to fail, thereby preventing the server from starting up, even when the problem was quite harmless. An example is that an unwritable file in the data directory would prevent restart on some platforms; but it is common to make SSL certificate files unwritable by the server. Revise this behavior so that permissions failures are ignored altogether, and other types of failures are logged but do not prevent continuing.

- Remove configure's check prohibiting linking to a threaded libpython on OpenBSD (Tom Lane)

The failure this restriction was meant to prevent seems to not be a problem anymore on current OpenBSD versions.

- Allow libpq to use TLS protocol versions beyond v1 (Noah Misch)

For a long time, libpq was coded so that the only SSL protocol it would allow was TLS v1. Now that newer TLS versions are becoming popular, allow it to negotiate the highest commonly-supported TLS version with the server. (PostgreSQL servers were already capable of such negotiation, so no change is needed on the server side.) This is a back-patch of a change already released in 9.4.0.

E.118. Release 9.1.16

Release date: 2015-05-22

This release contains a variety of fixes from 9.1.15. For information about new features in the 9.1 major release, see Section E.134.

E.118.1. Migration to Version 9.1.16

A dump/restore is not required for those running 9.1.X.

However, if you use contrib/citext's regexp_matches() functions, see the changelog entry below about that.

Also, if you are upgrading from a version earlier than 9.1.14, see Section E.120.

E.118.2. Changes

- Avoid possible crash when client disconnects just before the authentication timeout expires (Benkocs Norbert Attila)

If the timeout interrupt fired partway through the session shutdown sequence, SSL-related state would be freed twice, typically causing a crash and hence denial of service to other sessions. Experimentation shows that an unauthenticated remote attacker could trigger the bug somewhat consistently, hence treat as security issue. (CVE-2015-3165)

- Improve detection of system-call failures (Noah Misch)

Our replacement implementation of snprintf() failed to check for errors reported by the underlying system library calls; the main case that might be missed is out-of-memory situations. In the worst case this might lead to information exposure, due to our code assuming that a buffer had been overwritten when it hadn't been. Also, there were a few places in which security-relevant calls of other system library functions did not check for failure.
It remains possible that some calls of the *printf() family of functions are vulnerable to information disclosure if an out-of-memory error occurs at just the wrong time. We judge the risk to not be large, but will continue analysis in this area. (CVE-2015-3166)

- In contrib/pgcrypto, uniformly report decryption failures as “Wrong key or corrupt data” (Noah Misch)

Previously, some cases of decryption with an incorrect key could report other error message texts. It has been shown that such variance in error reports can aid attackers in recovering keys from other systems. While it's unknown whether pgcrypto's specific behaviors are likewise exploitable, it seems better to avoid the risk by using a one-size-fits-all message. (CVE-2015-3167)

- Fix incorrect declaration of contrib/citext's regexp_matches() functions (Tom Lane)

These functions should return setof text[], like the core functions they are wrappers for; but they were incorrectly declared as returning just text[]. This mistake had two results: first, if there was no match you got a scalar null result, whereas what you should get is an empty set (zero rows). Second, the q flag was effectively ignored, since you would get only one result array even if there were multiple matches.

While the latter behavior is clearly a bug, there might be applications depending on the former behavior; therefore the function declarations will not be changed by default until PostgreSQL 9.5. In pre-9.5 branches, the old behavior exists in version 1.0 of the citext extension, while we have provided corrected declarations in version 1.1 (which is not installed by default). To adopt the fix in pre-9.5 branches, execute ALTER EXTENSION citext UPDATE TO '1.1' in each database in which citext is installed. (You can also “update” back to 1.0 if you need to undo that.) Be aware that either update direction will require dropping and recreating any views or rules that use citext's regexp_matches() functions.

- Fix incorrect checking of deferred exclusion constraints after a HOT update (Tom Lane)

If a new row that potentially violates a deferred exclusion constraint is HOT-updated (that is, no indexed columns change and the row can be stored back onto the same table page) later in the same transaction, the exclusion constraint would be reported as violated when the check finally occurred, even if the row(s) the new row originally conflicted with had been deleted.

- Prevent improper reordering of antijoins (NOT EXISTS joins) versus other outer joins (Tom Lane)

This oversight in the planner has been observed to cause “could not find RelOptInfo for given relids” errors, but it seems possible that sometimes an incorrect query plan might get past that consistency check and result in silently-wrong query output.

- Fix incorrect matching of subexpressions in outer-join plan nodes (Tom Lane)

Previously, if textually identical non-strict subexpressions were used both above and below an outer join, the planner might try to re-use the value computed below the join, which would be incorrect because the executor would force the value to NULL in case of an unmatched outer row.

- Fix GEQO planner to cope with failure of its join order heuristic (Tom Lane)

This oversight has been seen to lead to “failed to join all relations together” errors in queries involving LATERAL, and that might happen in other cases as well.

- Fix possible deadlock at startup when max_prepared_transactions is too small (Heikki Linnakangas)

- Don't archive useless preallocated WAL files after a timeline switch (Heikki Linnakangas)

- Avoid “cannot GetMultiXactIdMembers() during recovery” error (Álvaro Herrera)

- Recursively fsync() the data directory after a crash (Abhijit Menon-Sen, Robert Haas)
This ensures consistency if another crash occurs shortly later. (The second crash would have to be a system-level crash, not just a database crash, for there to be a problem.)

- Fix autovacuum launcher’s possible failure to shut down, if an error occurs after it receives SIGTERM (Álvaro Herrera)

- Cope with unexpected signals in LockBufferForCleanup() (Andres Freund)

  This oversight could result in spurious errors about “multiple backends attempting to wait for pin-count 1”.

- Avoid waiting for WAL flush or synchronous replication during commit of a transaction that was read-only so far as the user is concerned (Andres Freund)

Previous, a delay could occur at commit in transactions that had written WAL due to HOT page pruning, leading to undesirable effects such as sessions getting stuck at startup if all synchronous replicas are down. Sessions have also been observed to get stuck in catchup interrupt processing when using synchronous replication; this will fix that problem as well.

- Fix crash when manipulating hash indexes on temporary tables (Heikki Linnakangas)

- Fix possible failure during hash index bucket split, if other processes are modifying the index concurrently (Tom Lane)

- Check for interrupts while analyzing index expressions (Jeff Janes)

  `ANALYZE` executes index expressions many times; if there are slow functions in such an expression, it’s desirable to be able to cancel the `ANALYZE` before that loop finishes.

- Ensure tableoid of a foreign table is reported correctly when a `READ COMMITTED` recheck occurs after locking rows in `SELECT FOR UPDATE`, `UPDATE`, or `DELETE` (Etsuro Fujita)

- Add the name of the target server to object description strings for foreign-server user mappings (Álvaro Herrera)

- Recommend setting `include_realm` to 1 when using Kerberos/GSSAPI/SSPI authentication (Stephen Frost)

  Without this, identically-named users from different realms cannot be distinguished. For the moment this is only a documentation change, but it will become the default setting in PostgreSQL 9.5.

- Remove code for matching IPv4 `pg_hba.conf` entries to IPv4-in-IPv6 addresses (Tom Lane)

  This hack was added in 2003 in response to a report that some Linux kernels of the time would report IPv4 connections as having IPv4-in-IPv6 addresses. However, the logic was accidentally broken in 9.0. The lack of any field complaints since then shows that it's not needed anymore. Now we have reports that the broken code causes crashes on some systems, so let's just remove it rather than fix it. (Had we chosen to fix it, that would make for a subtle and potentially security-sensitive change in the effective meaning of IPv4 `pg_hba.conf` entries, which does not seem like a good thing to do in minor releases.)

- Report WAL flush, not insert, position in `IDENTIFY_SYSTEM` replication command (Heikki Linnakangas)

  This avoids a possible startup failure in `pg_receivexlog`.

- While shutting down service on Windows, periodically send status updates to the Service Control Manager to prevent it from killing the service too soon; and ensure that `pg_ctl` will wait for shutdown (Krystian Bigaj)

- Reduce risk of network deadlock when using libpq’s non-blocking mode (Heikki Linnakangas)
When sending large volumes of data, it's important to drain the input buffer every so often, in case the server has sent enough response data to cause it to block on output. (A typical scenario is that the server is sending a stream of NOTICE messages during COPY FROM STDIN.) This worked properly in the normal blocking mode, but not so much in non-blocking mode. We've modified libpq to opportunistically drain input when it can, but a full defense against this problem requires application cooperation: the application should watch for socket read-ready as well as write-ready conditions, and be sure to call PQconsumeInput() upon read-ready.

- Fix array handling in ecpg (Michael Meskes)
- Fix psql to sanely handle URIs and conninfo strings as the first parameter to \connect (David Fetter, Andrew Dunstan, Álvaro Herrera)

This syntax has been accepted (but undocumented) for a long time, but previously some parameters might be taken from the old connection instead of the given string, which was agreed to be undesirable.

- Suppress incorrect complaints from psql on some platforms that it failed to write ~/.psql_history at exit (Tom Lane)

This misbehavior was caused by a workaround for a bug in very old (pre-2006) versions of libedit. We fixed it by removing the workaround, which will cause a similar failure to appear for anyone still using such versions of libedit. Recommendation: upgrade that library, or use libreadline.

- Fix pg_dump's rule for deciding which casts are system-provided casts that should not be dumped (Tom Lane)
- In pg_dump, fix failure to honor -Z compression level option together with -Fd (Michael Paquier)
- Make pg_dump consider foreign key relationships between extension configuration tables while choosing dump order (Gilles Darold, Michael Paquier, Stephen Frost)

This oversight could result in producing dumps that fail to reload because foreign key constraints are transiently violated.

- Fix dumping of views that are just VALUES(...) but have column aliases (Tom Lane)
- In pg_upgrade, force timeline 1 in the new cluster (Bruce Momjian)

This change prevents upgrade failures caused by bogus complaints about missing WAL history files.

- In pg_upgrade, check for improperly non-connectable databases before proceeding (Bruce Momjian)
- In pg_upgrade, quote directory paths properly in the generated delete_old_cluster script (Bruce Momjian)
- In pg_upgrade, preserve database-level freezing info properly (Bruce Momjian)

This oversight could cause missing-clog-file errors for tables within the postgres and template1 databases.

- Run pg_upgrade and pg_resetzxlog with restricted privileges on Windows, so that they don't fail when run by an administrator (Muhammad Asif Naeem)
- Improve handling of readdir() failures when scanning directories in initdb and pg_basebackup (Marco Nenciarini)
- Fix slow sorting algorithm in contrib/intarray (Tom Lane)
- Fix compile failure on Sparc V8 machines (Rob Rowan)
Release Notes

- Update time zone data files to tzdata release 2015d for DST law changes in Egypt, Mongolia, and Palestine, plus historical changes in Canada and Chile. Also adopt revised zone abbreviations for the America/Adak zone (HST/HDT not HAST/HADT).

E.119. Release 9.1.15

Release date: 2015-02-05

This release contains a variety of fixes from 9.1.14. For information about new features in the 9.1 major release, see Section E.134.

E.119.1. Migration to Version 9.1.15

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.14, see Section E.120.

E.119.2. Changes

- Fix buffer overruns in `to_char()` (Bruce Momjian)

  When `to_char()` processes a numeric formatting template calling for a large number of digits, PostgreSQL would read past the end of a buffer. When processing a crafted timestamp formatting template, PostgreSQL would write past the end of a buffer. Either case could crash the server. We have not ruled out the possibility of attacks that lead to privilege escalation, though they seem unlikely. (CVE-2015-0241)

- Fix buffer overrun in replacement `printf()` functions (Tom Lane)

  PostgreSQL includes a replacement implementation of `printf` and related functions. This code will overrun a stack buffer when formatting a floating point number (conversion specifiers e, E, f, F, g or G) with requested precision greater than about 500. This will crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. A database user can trigger such a buffer overrun through the `to_char()` SQL function. While that is the only affected core PostgreSQL functionality, extension modules that use printf-family functions may be at risk as well.

  This issue primarily affects PostgreSQL on Windows. PostgreSQL uses the system implementation of these functions where adequate, which it is on other modern platforms. (CVE-2015-0242)

- Fix buffer overruns in `contrib/pgcrypto` (Marko Tiikkaja, Noah Misch)

  Errors in memory size tracking within the `pgcrypto` module permitted stack buffer overruns and improper dependence on the contents of uninitialized memory. The buffer overrun cases can crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. (CVE-2015-0243)

- Fix possible loss of frontend/backend protocol synchronization after an error (Heikki Linnakangas)

  If any error occurred while the server was in the middle of reading a protocol message from the client, it could lose synchronization and incorrectly try to interpret part of the message's data as a new protocol message. An attacker able to submit crafted binary data within a command parameter might succeed in injecting his own SQL commands this way. Statement timeout and query cancellation are the most likely sources of errors triggering this scenario. Particularly vulnerable are applications that use a timeout and also submit arbitrary user-crafted data as binary query parameters. Disabling statement timeout will reduce, but not eliminate, the risk of exploit. Our thanks to Emil Lenngren for reporting this issue. (CVE-2015-0244)

- Fix information leak via constraint-violation error messages (Stephen Frost)
Some server error messages show the values of columns that violate a constraint, such as a unique constraint. If the user does not have SELECT privilege on all columns of the table, this could mean exposing values that the user should not be able to see. Adjust the code so that values are displayed only when they came from the SQL command or could be selected by the user. (CVE-2014-8161)

- Lock down regression testing's temporary installations on Windows (Noah Misch)

  Use SSPI authentication to allow connections only from the OS user who launched the test suite. This closes on Windows the same vulnerability previously closed on other platforms, namely that other users might be able to connect to the test postmaster. (CVE-2014-0067)

- Avoid possible data corruption if `ALTER DATABASE SET TABLESPACE` is used to move a database to a new tablespace and then shortly later move it back to its original tablespace (Tom Lane)

- Avoid corrupting tables when `ANALYZE` inside a transaction is rolled back (Andres Freund, Tom Lane, Michael Paquier)

  If the failing transaction had earlier removed the last index, rule, or trigger from the table, the table would be left in a corrupted state with the relevant `pg_class` flags not set though they should be.

- Ensure that unlogged tables are copied correctly during `CREATE DATABASE` or `ALTER DATABASE SET TABLESPACE` (Pavan Deolasee, Andres Freund)

- Fix `DROP`’s dependency searching to correctly handle the case where a table column is recursively visited before its table (Petr Jelinek, Tom Lane)

  This case is only known to arise when an extension creates both a datatype and a table using that datatype. The faulty code might refuse a `DROP EXTENSION` unless `CASCADE` is specified, which should not be required.

- Fix use-of-already-freed-memory problem in EvalPlanQual processing (Tom Lane)

  In `READ COMMITTED` mode, queries that lock or update recently-updated rows could crash as a result of this bug.

- Fix planning of `SELECT FOR UPDATE` when using a partial index on a child table (Kyotaro Horiguchi)

  In `READ COMMITTED` mode, `SELECT FOR UPDATE` must also recheck the partial index’s `WHERE` condition when rechecking a recently-updated row to see if it still satisfies the query’s `WHERE` condition. This requirement was missed if the index belonged to an inheritance child table, so that it was possible to incorrectly return rows that no longer satisfy the query condition.

- Fix corner case wherein `SELECT FOR UPDATE` could return a row twice, and possibly miss returning other rows (Tom Lane)

  In `READ COMMITTED` mode, a `SELECT FOR UPDATE` that is scanning an inheritance tree could incorrectly return a row from a prior child table instead of the one it should return from a later child table.

- Reject duplicate column names in the referenced-columns list of a `FOREIGN KEY` declaration (David Rowley)

  This restriction is per SQL standard. Previously we did not reject the case explicitly, but later on the code would fail with bizarre-looking errors.

- Fix bugs in raising a numeric value to a large integral power (Tom Lane)

  The previous code could get a wrong answer, or consume excessive amounts of time and memory before realizing that the answer must overflow.
• In `numeric_recv()`, truncate away any fractional digits that would be hidden according to the value's `dscale` field (Tom Lane)

A numeric value's display scale (`dscale`) should never be less than the number of nonzero fractional digits; but apparently there's at least one broken client application that transmits binary numeric values in which that's true. This leads to strange behavior since the extra digits are taken into account by arithmetic operations even though they aren't printed. The least risky fix seems to be to truncate away such “hidden” digits on receipt, so that the value is indeed what it prints as.

• Reject out-of-range numeric timezone specifications (Tom Lane)

Simple numeric timezone specifications exceeding +/- 168 hours (one week) would be accepted, but could then cause null-pointer dereference crashes in certain operations. There's no use-case for such large UTC offsets, so reject them.

• Fix bugs in `tsquery @> tsquery` operator (Heikki Linnakangas)

Two different terms would be considered to match if they had the same CRC. Also, if the second operand had more terms than the first, it would be assumed not to be contained in the first; which is wrong since it might contain duplicate terms.

• Improve `ispell` dictionary's defenses against bad affix files (Tom Lane)

• Allow more than 64K phrases in a thesaurus dictionary (David Boutin)

The previous coding could crash on an oversize dictionary, so this was deemed a back-patchable bug fix rather than a feature addition.

• Fix namespace handling in `xpath()` (Ali Akbar)

Previously, the `xml` value resulting from an `xpath()` call would not have namespace declarations if the namespace declarations were attached to an ancestor element in the input `xml` value, rather than to the specific element being returned. Propagate the ancestral declaration so that the result is correct when considered in isolation.

• Fix planner problems with nested append relations, such as inherited tables within `UNION ALL` subqueries (Tom Lane)

• Fail cleanly when a GiST index tuple doesn't fit on a page, rather than going into infinite recursion (Andrew Gierth)

• Exempt tables that have per-table `cost_limit` and/or `cost_delay` settings from autovacuum's global cost balancing rules (Álvaro Herrera)

The previous behavior resulted in basically ignoring these per-table settings, which was unintended. Now, a table having such settings will be vacuumed using those settings, independently of what is going on in other autovacuum workers. This may result in heavier total I/O load than before, so such settings should be re-examined for sanity.

• Avoid wholesale autovacuuming when autovacuum is nominally off (Tom Lane)

Even when autovacuum is nominally off, we will still launch autovacuum worker processes to vacuum tables that are at risk of XID wraparound. However, such a worker process then proceeded to vacuum all tables in the target database, if they met the usual thresholds for autovacuuming. This is at best pretty unexpected; at worst it delays response to the wraparound threat. Fix it so that if autovacuum is turned off, workers only do anti-wraparound vacuums and not any other work.

• During crash recovery, ensure that unlogged relations are rewritten as empty and are synced to disk before recovery is considered complete (Abhijit Menon-Sen, Andres Freund)

This prevents scenarios in which unlogged relations might contain garbage data following database crash recovery.
• Fix race condition between hot standby queries and replaying a full-page image (Heikki Linnakan-
gas)

This mistake could result in transient errors in queries being executed in hot standby.

• Fix several cases where recovery logic improperly ignored WAL records for COMMIT/ABORT
PREPARED (Heikki Linnakangas)

The most notable oversight was that recovery_target_xid could not be used to stop at a two-
phase commit.

• Avoid creating unnecessary .ready marker files for timeline history files (Fujii Masao)

• Fix possible null pointer dereference when an empty prepared statement is used and the log_s-
tatement setting is mod or ddl (Fujii Masao)

• Change “pgstat wait timeout” warning message to be LOG level, and rephrase it to be more under-
standable (Tom Lane)

This message was originally thought to be essentially a can’t-happen case, but it occurs often enough
on our slower buildfarm members to be a nuisance. Reduce it to LOG level, and expend a bit more
effort on the wording: it now reads "using stale statistics instead of current ones because stats col-
lector is not responding”.

• Fix SPARC spinlock implementation to ensure correctness if the CPU is being run in a non-TSO
coherency mode, as some non-Solaris kernels do (Andres Freund)

• Warn if macOS’s setlocale() starts an unwanted extra thread inside the postmaster (Noah
Misch)

• Fix processing of repeated dbname parameters in PQconnectdbParams() (Alex Shulgin)

Unexpected behavior ensued if the first occurrence of dbname contained a connection string or
URI to be expanded.

• Ensure that libpq reports a suitable error message on unexpected socket EOF (Marko Tiikkaja, Tom
Lane)

Depending on kernel behavior, libpq might return an empty error string rather than something useful
when the server unexpectedly closed the socket.

• Clear any old error message during PQreset () (Heikki Linnakangas)

If PQreset() is called repeatedly, and the connection cannot be re-established, error messages
from the failed connection attempts kept accumulating in the PGconn’s error string.

• Properly handle out-of-memory conditions while parsing connection options in libpq (Alex Shulgin,
Heikki Linnakangas)

• Fix array overrun in ecpg’s version of ParseDateTime() (Michael Paquier)

• In initdb, give a clearer error message if a password file is specified but is empty (Mats Erik An-
dersson)

• Fix psql’s \s command to work nicely with libedit, and add pager support (Stepan Rutz, Tom Lane)

When using libedit rather than readline, \s printed the command history in a fairly unreadable
encoded format, and on recent libedit versions might fail altogether. Fix that by printing the history
ourselves rather than having the library do it. A pleasant side-effect is that the pager is used if
appropriate.

This patch also fixes a bug that caused newline encoding to be applied inconsistently when saving
the command history with libedit. Multiline history entries written by older psql versions will be
read cleanly with this patch, but perhaps not vice versa, depending on the exact libedit versions involved.

- Improve consistency of parsing of psql's special variables (Tom Lane)

  Allow variant spellings of on and off (such as 1/0) for ECHO_HIDDEN and ON_ERROR_ROLLBACK. Report a warning for unrecognized values for COMP_KEYWORD_CASE, ECHO, ECHO_HIDDEN, HISTCONTROL, ON_ERROR_ROLLBACK, and VERBOSITY. Recognize all values for all these variables case-insensitively; previously there was a mishmash of case-sensitive and case-insensitive behaviors.

- Fix psql's expanded-mode display to work consistently when using border = 3 and linestyle = ascii or unicode (Stephen Frost)

- Improve performance of pg_dump when the database contains many instances of multiple dependency paths between the same two objects (Tom Lane)

- Fix possible deadlock during parallel restore of a schema-only dump (Robert Haas, Tom Lane)

- Fix core dump in pg_dump --binary-upgrade on zero-column composite type (Rushabh Lathia)

- Prevent WAL files created by pg_basebackup -x/-X from being archived again when the standby is promoted (Andres Freund)

- Fix upgrade-from-unpackaged script for contrib/citext (Tom Lane)

- Fix block number checking in contrib/pageinspect's get_raw_page() (Tom Lane)

  The incorrect checking logic could prevent access to some pages in non-main relation forks.

- Fix contrib/pgcrypto's pgp_sym_decrypt() to not fail on messages whose length is 6 less than a power of 2 (Marko Tiikkaja)

- Fix file descriptor leak in contrib/pg_test_fsync (Jeff Janes)

  This could cause failure to remove temporary files on Windows.

- Handle unexpected query results, especially NULLs, safely in contrib/tablefunc's connectby() (Michael Paquier)

  connectby() previously crashed if it encountered a NULL key value. It now prints that row but doesn't recurse further.

- Avoid a possible crash in contrib/xml2's xslt_process() (Mark Simonetti)

  libxml seems to have an undocumented dependency on the order in which resources are freed; reorder our calls to avoid a crash.

- Mark some contrib I/O functions with correct volatility properties (Tom Lane)

  The previous over-conservative marking was immaterial in normal use, but could cause optimization problems or rejection of valid index expression definitions. Since the consequences are not large, we've just adjusted the function definitions in the extension modules' scripts, without changing version numbers.

- Numerous cleanups of warnings from Coverity static code analyzer (Andres Freund, Tatsuo Ishii, Marko Kreen, Tom Lane, Michael Paquier)

  These changes are mostly cosmetic but in some cases fix corner-case bugs, for example a crash rather than a proper error report after an out-of-memory failure. None are believed to represent security issues.
• Detect incompatible OpenLDAP versions during build (Noah Misch)

With OpenLDAP versions 2.4.24 through 2.4.31, inclusive, PostgreSQL backends can crash at exit. Raise a warning during configure based on the compile-time OpenLDAP version number, and test the crashing scenario in the contrib/dblink regression test.

• In non-MSVC Windows builds, ensure libpq.dll is installed with execute permissions (Noah Misch)

• Make pg_regress remove any temporary installation it created upon successful exit (Tom Lane)

This results in a very substantial reduction in disk space usage during make check-world, since that sequence involves creation of numerous temporary installations.

• Support time zone abbreviations that change UTC offset from time to time (Tom Lane)

Previously, PostgreSQL assumed that the UTC offset associated with a time zone abbreviation (such as EST) never changes in the usage of any particular locale. However this assumption fails in the real world, so introduce the ability for a zone abbreviation to represent a UTC offset that sometimes changes. Update the zone abbreviation definition files to make use of this feature in timezone locales that have changed the UTC offset of their abbreviations since 1970 (according to the IANA timezone database). In such timezones, PostgreSQL will now associate the correct UTC offset with the abbreviation depending on the given date.

• Update time zone abbreviations lists (Tom Lane)

Add CST (China Standard Time) to our lists. Remove references to ADT as “Arabia Daylight Time”, an abbreviation that’s been out of use since 2007; therefore, claiming there is a conflict with “Atlantic Daylight Time” doesn’t seem especially helpful. Fix entirely incorrect GMT offsets for CKT (Cook Islands), FJT, and FJST (Fiji); we didn’t even have them on the proper side of the date line.

• Update time zone data files to tzdata release 2015a.

The IANA timezone database has adopted abbreviations of the form A\times T/A\times DT for all Australian time zones, reflecting what they believe to be current majority practice Down Under. These names do not conflict with usage elsewhere (other than ACST for Acre Summer Time, which has been in disuse since 1994). Accordingly, adopt these names into our “Default” timezone abbreviation set. The “Australia” abbreviation set now contains only CST, EAST, EST, SAST, SAT, and WST, all of which are thought to be mostly historical usage. Note that SAST has also been changed to be South Africa Standard Time in the “Default” abbreviation set.

Also, add zone abbreviations SRET (Asia/Srednekolymsk) and XJT (Asia/Urumqi), and use WSST/ WSDT for western Samoa. Also, there were DST law changes in Chile, Mexico, the Turks & Caicos Islands (America/Grand_Turk), and Fiji. There is a new zone Pacific/Bougainville for portions of Papua New Guinea. Also, numerous corrections for historical (pre-1970) time zone data.

E.120. Release 9.1.14

Release date: 2014-07-24

This release contains a variety of fixes from 9.1.13. For information about new features in the 9.1 major release, see Section E.134.

E.120.1. Migration to Version 9.1.14

A dump/restore is not required for those running 9.1.X.

However, this release corrects an index corruption problem in some GiST indexes. See the first changelog entry below to find out whether your installation has been affected and what steps you should take if so.
Also, if you are upgrading from a version earlier than 9.1.11, see Section E.123.

E.120.2. Changes

• Correctly initialize padding bytes in contrib/btree_gist indexes on bit columns (Heikki Linnakangas)
  This error could result in incorrect query results due to values that should compare equal not being seen as equal. Users with GiST indexes on bit or bit varying columns should REINDEX those indexes after installing this update.

• Protect against torn pages when deleting GIN list pages (Heikki Linnakangas)
  This fix prevents possible index corruption if a system crash occurs while the page update is being written to disk.

• Don’t clear the right-link of a GiST index page while replaying updates from WAL (Heikki Linnakangas)
  This error could lead to transiently wrong answers from GiST index scans performed in Hot Standby.

• Fix feedback status when hot_standby_feedback is turned off on-the-fly (Simon Riggs)

• Fix possibly-incorrect cache invalidation during nested calls to ReceiveSharedInvalidMessages (Andres Freund)

• Fix “could not find pathkey item to sort” planner failures with UNION ALL over subqueries reading from tables with inheritance children (Tom Lane)

• Don’t assume a subquery’s output is unique if there’s a set-returning function in its targetlist (David Rowley)
  This oversight could lead to misoptimization of constructs like WHERE x IN (SELECT y, generate_series(1,10) FROM t GROUP BY y).

• Fix failure to detoast fields in composite elements of structured types (Tom Lane)
  This corrects cases where TOAST pointers could be copied into other tables without being dereferenced. If the original data is later deleted, it would lead to errors like “missing chunk number 0 for toast value ...” when the now-dangling pointer is used.

• Fix “record type has not been registered” failures with whole-row references to the output of Append plan nodes (Tom Lane)

• Fix possible crash when invoking a user-defined function while rewinding a cursor (Tom Lane)

• Fix query-lifespan memory leak while evaluating the arguments for a function in FROM (Tom Lane)

• Fix session-lifespan memory leaks in regular-expression processing (Tom Lane, Arthur O'Dwyer, Greg Stark)

• Fix data encoding error in hungarian.stop (Tom Lane)

• Prevent foreign tables from being created with OIDS when default_with_oids is true (Etsuro Fujita)

• Fix liveness checks for rows that were inserted in the current transaction and then deleted by a now-rolled-back subtransaction (Andres Freund)
  This could cause problems (at least spurious warnings, and at worst an infinite loop) if CREATE INDEX or CLUSTER were done later in the same transaction.

• Clear pg_stat_activity.xact_start during PREPARE TRANSACTION (Andres Freund)
After the PREPARE, the originating session is no longer in a transaction, so it should not continue to display a transaction start time.

- **Fix REASSIGN OWNED** to not fail for text search objects (Álvaro Herrera)
- **Block signals during postmaster startup** (Tom Lane)

  This ensures that the postmaster will properly clean up after itself if, for example, it receives SIGINT while still starting up.

- **Fix client host name lookup when processing pg_hba.conf entries that specify host names instead of IP addresses** (Tom Lane)

  Ensure that reverse-DNS lookup failures are reported, instead of just silently not matching such entries. Also ensure that we make only one reverse-DNS lookup attempt per connection, not one per host name entry, which is what previously happened if the lookup attempts failed.

- **Secure Unix-domain sockets of temporary postmasters started during make check** (Noah Misch)

  Any local user able to access the socket file could connect as the server's bootstrap superuser, then proceed to execute arbitrary code as the operating-system user running the test, as we previously noted in CVE-2014-0067. This change defends against that risk by placing the server's socket in a temporary, mode 0700 subdirectory of /tmp. The hazard remains however on platforms where Unix sockets are not supported, notably Windows, because then the temporary postmaster must accept local TCP connections.

  A useful side effect of this change is to simplify make check testing in builds that override DEFAULT_PGSOCKET_DIR. Popular non-default values like /var/run/postgresql are often not writable by the build user, requiring workarounds that will no longer be necessary.

- **Fix tablespace creation WAL replay to work on Windows** (MauMau)
- **Fix detection of socket creation failures on Windows** (Bruce Momjian)
- **On Windows, allow new sessions to absorb values of PGC_BACKEND parameters** (such as log_connections) from the configuration file (Amit Kapila)

  Previously, if such a parameter were changed in the file post-startup, the change would have no effect.

- **Properly quote executable path names on Windows** (Nikhil Deshpande)

  This oversight could cause initdb and pg_upgrade to fail on Windows, if the installation path contained both spaces and @ signs.

- **Fix linking of libpython on macOS** (Tom Lane)

  The method we previously used can fail with the Python library supplied by Xcode 5.0 and later.

- **Avoid buffer bloat in libpq when the server consistently sends data faster than the client can absorb it** (Shin-ichi Morita, Tom Lane)

  libpq could be coerced into enlarging its input buffer until it runs out of memory (which would be reported misleadingly as “lost synchronization with server”). Under ordinary circumstances it's quite far-fetched that data could be continuously transmitted more quickly than the recv() loop can absorb it, but this has been observed when the client is artificially slowed by scheduler constraints.

- **Ensure that LDAP lookup attempts in libpq time out as intended** (Laurenz Albe)

- **Fix ecpg to do the right thing when an array of char * is the target for a FETCH statement returning more than one row, as well as some other array-handling fixes** (Ashutosh Bapat)
• Fix pg_restore's processing of old-style large object comments (Tom Lane)

A direct-to-database restore from an archive file generated by a pre-9.0 version of pg_dump would usually fail if the archive contained more than a few comments for large objects.

• In contrib/pgcrypto functions, ensure sensitive information is cleared from stack variables before returning (Marko Kreen)

• In contrib/uuid-ossp, cache the state of the OSSP UUID library across calls (Tom Lane)

This improves the efficiency of UUID generation and reduces the amount of entropy drawn from /dev/urandom, on platforms that have that.

• Update time zone data files to tzdata release 2014e for DST law changes in Crimea, Egypt, and Morocco.

E.121. Release 9.1.13

Release date: 2014-03-20

This release contains a variety of fixes from 9.1.12. For information about new features in the 9.1 major release, see Section E.134.

E.121.1. Migration to Version 9.1.13

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.11, see Section E.123.

E.121.2. Changes

• Restore GIN metapages unconditionally to avoid torn-page risk (Heikki Linnakangas)

Although this oversight could theoretically result in a corrupted index, it is unlikely to have caused any problems in practice, since the active part of a GIN metapage is smaller than a standard 512-byte disk sector.

• Avoid race condition in checking transaction commit status during receipt of a NOTIFY message (Marko Tiikkaja)

This prevents a scenario wherein a sufficiently fast client might respond to a notification before database updates made by the notifier have become visible to the recipient.

• Allow regular-expression operators to be terminated early by query cancel requests (Tom Lane)

This prevents scenarios wherein a pathological regular expression could lock up a server process uninterruptibly for a long time.

• Remove incorrect code that tried to allow OVERLAPS with single-element row arguments (Joshua Yanovski)

This code never worked correctly, and since the case is neither specified by the SQL standard nor documented, it seemed better to remove it than fix it.

• Avoid getting more than AccessShareLock when de-parsing a rule or view (Dean Rasheed)

This oversight resulted in pg_dump unexpectedly acquiring RowExclusiveLock locks on tables mentioned as the targets of INSERT/UPDATE/DELETE commands in rules. While usually harm-
less, that could interfere with concurrent transactions that tried to acquire, for example, Share-Lock on those tables.

- Improve performance of index endpoint probes during planning (Tom Lane)

  This change fixes a significant performance problem that occurred when there were many not-yet-committed rows at the end of the index, which is a common situation for indexes on sequentially-assigned values such as timestamps or sequence-generated identifiers.

- Fix walsender’s failure to shut down cleanly when client is pg_receivexlog (Fujii Masao)

- Fix test to see if hot standby connections can be allowed immediately after a crash (Heikki Linnakangas)

- Prevent interrupts while reporting non-ERROR messages (Tom Lane)

  This guards against rare server-process freezeups due to recursive entry to syslog(), and perhaps other related problems.

- Fix memory leak in PL/Perl when returning a composite result, including multiple-OUT-parameter cases (Alex Hunsaker)

- Prevent intermittent “could not reserve shared memory region” failures on recent Windows versions (MauMau)

- Update time zone data files to tzdata release 2014a for DST law changes in Fiji and Turkey, plus historical changes in Israel and Ukraine.

**E.122. Release 9.1.12**

**Release date:** 2014-02-20

This release contains a variety of fixes from 9.1.11. For information about new features in the 9.1 major release, see Section E.134.

**E.122.1. Migration to Version 9.1.12**

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.11, see Section E.123.

**E.122.2. Changes**

- Shore up GRANT ... WITH ADMIN OPTION restrictions (Noah Misch)

  Granting a role without ADMIN OPTION is supposed to prevent the grantee from adding or removing members from the granted role, but this restriction was easily bypassed by doing SET ROLE first. The security impact is mostly that a role member can revoke the access of others, contrary to the wishes of his grantor. Unapproved role member additions are a lesser concern, since an uncooperative role member could provide most of his rights to others anyway by creating views or SECURITY DEFINER functions. (CVE-2014-0060)

- Prevent privilege escalation via manual calls to PL validator functions (Andres Freund)

  The primary role of PL validator functions is to be called implicitly during CREATE FUNCTION, but they are also normal SQL functions that a user can call explicitly. Calling a validator on a function actually written in some other language was not checked for and could be exploited for privilege-escalation purposes. The fix involves adding a call to a privilege-checking function in each validator function. Non-core procedural languages will also need to make this change to their own validator functions, if any. (CVE-2014-0061)
• Avoid multiple name lookups during table and index DDL (Robert Haas, Andres Freund)

If the name lookups come to different conclusions due to concurrent activity, we might perform some parts of the DDL on a different table than other parts. At least in the case of CREATE INDEX, this can be used to cause the permissions checks to be performed against a different table than the index creation, allowing for a privilege escalation attack. (CVE-2014-0062)

• Prevent buffer overrun with long datetime strings (Noah Misch)

The MAXDATELEN constant was too small for the longest possible value of type interval, allowing a buffer overrun in interval_out(). Although the datetime input functions were more careful about avoiding buffer overrun, the limit was short enough to cause them to reject some valid inputs, such as input containing a very long timezone name. The ecpg library contained these vulnerabilities along with some of its own. (CVE-2014-0063)

• Prevent buffer overrun due to integer overflow in size calculations (Noah Misch, Heikki Linnakan-gas)

Several functions, mostly type input functions, calculated an allocation size without checking for overflow. If overflow did occur, a too-small buffer would be allocated and then written past. (CVE-2014-0064)

• Prevent overruns of fixed-size buffers (Peter Eisentraut, Jozef Mlich)

Use strncpy() and related functions to provide a clear guarantee that fixed-size buffers are not overrun. Unlike the preceding items, it is unclear whether these cases really represent live issues, since in most cases there appear to be previous constraints on the size of the input string. Nonetheless it seems prudent to silence all Coverity warnings of this type. (CVE-2014-0065)

• Avoid crashing if crypt() returns NULL (Honza Horak, Bruce Momjian)

There are relatively few scenarios in which crypt() could return NULL, but contrib/chk-pass would crash if it did. One practical case in which this could be an issue is if libc is configured to refuse to execute unapproved hashing algorithms (e.g., “FIPS mode”). (CVE-2014-0066)

• Document risks of make check in the regression testing instructions (Noah Misch, Tom Lane)

Since the temporary server started by make check uses “trust” authentication, another user on the same machine could connect to it as database superuser, and then potentially exploit the privileges of the operating-system user who started the tests. A future release will probably incorporate changes in the testing procedure to prevent this risk, but some public discussion is needed first. So for the moment, just warn people against using make check when there are untrusted users on the same machine. (CVE-2014-0067)

• Fix possible mis-replay of WAL records when some segments of a relation aren’t full size (Greg Stark, Tom Lane)

The WAL update could be applied to the wrong page, potentially many pages past where it should have been. Aside from corrupting data, this error has been observed to result in significant “bloat” of standby servers compared to their masters, due to updates being applied far beyond where the end-of-file should have been. This failure mode does not appear to be a significant risk during crash recovery, only when initially synchronizing a standby created from a base backup taken from a quickly-changing master.

• Fix bug in determining when recovery has reached consistency (Tomonari Katsumata, Heikki Lin-nakangas)

In some cases WAL replay would mistakenly conclude that the database was already consistent at the start of replay, thus possibly allowing hot-standby queries before the database was really consistent. Other symptoms such as “PANIC: WAL contains references to invalid pages” were also possible.
• Fix improper locking of btree index pages while replaying a VACUUM operation in hot-standby mode (Andres Freund, Heikki Linnakangas, Tom Lane)

This error could result in “PANIC: WAL contains references to invalid pages” failures.

• Ensure that insertions into non-leaf GIN index pages write a full-page WAL record when appropriate (Heikki Linnakangas)

The previous coding risked index corruption in the event of a partial-page write during a system crash.

• When pause_at_recovery_target and recovery_target_inclusive are both set, ensure the target record is applied before pausing, not after (Heikki Linnakangas)

• Fix race conditions during server process exit (Robert Haas)

Ensure that signal handlers don't attempt to use the process's MyProc pointer after it's no longer valid.

• Fix race conditions in walsender shutdown logic and walreceiver SIGHUP signal handler (Tom Lane)

• Fix unsafe references to errno within error reporting logic (Christian Kruse)

This would typically lead to odd behaviors such as missing or inappropriate HINT fields.

• Fix possible crashes from using ereport() too early during server startup (Tom Lane)

The principal case we've seen in the field is a crash if the server is started in a directory it doesn't have permission to read.

• Clear retry flags properly in OpenSSL socket write function (Alexander Kukushkin)

This omission could result in a server lockup after unexpected loss of an SSL-encrypted connection.

• Fix length checking for Unicode identifiers (U&"..." syntax) containing escapes (Tom Lane)

A spurious truncation warning would be printed for such identifiers if the escaped form of the identifier was too long, but the identifier actually didn't need truncation after de-escaping.

• Allow keywords that are type names to be used in lists of roles (Stephen Frost)

A previous patch allowed such keywords to be used without quoting in places such as role identifiers; but it missed cases where a list of role identifiers was permitted, such as DROP ROLE.

• Fix parser crash for EXISTS(SELECT * FROM zero_column_table) (Tom Lane)

• Fix possible crash due to invalid plan for nested sub-selects, such as WHERE (... x IN (SELECT ...) ...) IN (SELECT ...) (Tom Lane)

• Ensure that ANALYZE creates statistics for a table column even when all the values in it are “too wide” (Tom Lane)

ANALYZE intentionally omits very wide values from its histogram and most-common-values calculations, but it neglected to do something sane in the case that all the sampled entries are too wide.

• In ALTER TABLE ... SET TABLESPACE, allow the database's default tablespace to be used without a permissions check (Stephen Frost)

CREATE TABLE has always allowed such usage, but ALTER TABLE didn't get the memo.

• Fix “cannot accept a set” error when some arms of a CASE return a set and others don't (Tom Lane)
• Fix checks for all-zero client addresses in pgstat functions (Kevin Grittner)

• Fix possible misclassification of multibyte characters by the text search parser (Tom Lane)

Non-ASCII characters could be misclassified when using C locale with a multibyte encoding. On Cygwin, non-C locales could fail as well.

• Fix possible misbehavior in plainto_tsquery() (Heikki Linnakangas)

Use memmove() not memcpy() for copying overlapping memory regions. There have been no field reports of this actually causing trouble, but it's certainly risky.

• Fix placement of permissions checks in pg_start_backup() and pg_stop_backup() (Andres Freund, Magnus Hagander)

The previous coding might attempt to do catalog access when it shouldn't.

• Accept SHIFT_JIS as an encoding name for locale checking purposes (Tatsuo Ishii)

• Fix misbehavior of PQhost() on Windows (Fujii Masao)

It should return localhost if no host has been specified.

• Improve error handling in libpq and psql for failures during COPY TO STDOUT/FROM STDIN (Tom Lane)

In particular this fixes an infinite loop that could occur in 9.2 and up if the server connection was lost during COPY FROM STDIN. Variants of that scenario might be possible in older versions, or with other client applications.

• Fix possible incorrect printing of filenames in pg_basebackup's verbose mode (Magnus Hagander)

• Avoid including tablespaces inside PGDATA twice in base backups (Dimitri Fontaine, Magnus Hagander)

• Fix misaligned descriptors in ecpg (MauMau)

• In ecpg, handle lack of a hostname in the connection parameters properly (Michael Meskes)

• Fix performance regression in contrib/dblink connection startup (Joe Conway)

Avoid an unnecessary round trip when client and server encodings match.

• In contrib/ism, fix incorrect calculation of the check digit for ISMN values (Fabien Coelho)

• Ensure client-code-only installation procedure works as documented (Peter Eisentraut)

• In Mingw and Cygwin builds, install the libpq DLL in the bin directory (Andrew Dunstan)

This duplicates what the MSVC build has long done. It should fix problems with programs like psql failing to start because they can't find the DLL.

• Avoid using the deprecated dllwrap tool in Cygwin builds (Marco Atzeri)

• Don't generate plain-text HISTORY and src/test/regress/README files anymore (Tom Lane)

These text files duplicated the main HTML and PDF documentation formats. The trouble involved in maintaining them greatly outweighs the likely audience for plain-text format. Distribution tarballs will still contain files by these names, but they'll just be stubs directing the reader to consult the main documentation. The plain-text INSTALL file will still be maintained, as there is arguably a use-case for that.
• Update time zone data files to tzdata release 2013i for DST law changes in Jordan and historical changes in Cuba.

In addition, the zones Asia/Riyadh87, Asia/Riyadh88, and Asia/Riyadh89 have been removed, as they are no longer maintained by IANA, and never represented actual civil timekeeping practice.

### E.123. Release 9.1.11

**Release date:** 2013-12-05

This release contains a variety of fixes from 9.1.10. For information about new features in the 9.1 major release, see Section E.134.

### E.123.1. Migration to Version 9.1.11

A dump/restore is not required for those running 9.1.X. However, this release corrects a number of potential data corruption issues. See the first two changelog entries below to find out whether your installation has been affected and what steps you can take if so.

Also, if you are upgrading from a version earlier than 9.1.9, see Section E.125.

### E.123.2. Changes

- **Fix VACUUM's tests to see whether it can update relfrozenxid (Andres Freund)**

  In some cases VACUUM (either manual or autovacuum) could incorrectly advance a table's relfrozenxid value, allowing tuples to escape freezing, causing those rows to become invisible once 2^31 transactions have elapsed. The probability of data loss is fairly low since multiple incorrect advancements would need to happen before actual loss occurs, but it's not zero. Users upgrading from releases 9.0.4 or 8.4.8 or earlier are not affected, but all later versions contain the bug.

  The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix any latent corruption but will not be able to fix all pre-existing data errors. However, an installation can be presumed safe after performing this vacuuming if it has executed fewer than 2^31 update transactions in its lifetime (check this with SELECT txid_current() < 2^31).

- **Fix initialization of pg_clog and pg_subtrans during hot standby startup (Andres Freund, Heikki Linnakangas)**

  This bug can cause data loss on standby servers at the moment they start to accept hot-standby queries, by marking committed transactions as uncommitted. The likelihood of such corruption is small unless, at the time of standby startup, the primary server has executed many updating transactions since its last checkpoint. Symptoms include missing rows, rows that should have been deleted being still visible, and obsolete versions of updated rows being still visible alongside their newer versions.

  This bug was introduced in versions 9.3.0, 9.2.5, 9.1.10, and 9.0.14. Standby servers that have only been running earlier releases are not at risk. It's recommended that standby servers that have ever run any of the buggy releases be re-cloned from the primary (e.g., with a new base backup) after upgrading.

- **Truncate pg_multixact contents during WAL replay (Andres Freund)**

  This avoids ever-increasing disk space consumption in standby servers.

- **Fix race condition in GIN index posting tree page deletion (Heikki Linnakangas)**
This could lead to transient wrong answers or query failures.

- Avoid flattening a subquery whose SELECT list contains a volatile function wrapped inside a sub-SELECT (Tom Lane)

  This avoids unexpected results due to extra evaluations of the volatile function.

- Fix planner's processing of non-simple-variable subquery outputs nested within outer joins (Tom Lane)

  This error could lead to incorrect plans for queries involving multiple levels of subqueries within JOIN syntax.

- Fix incorrect generation of optimized MIN() / MAX() plans for inheritance trees (Tom Lane)

  The planner could fail in cases where the MIN() / MAX() argument was an expression rather than a simple variable.

- Fix premature deletion of temporary files (Andres Freund)

- Fix possible read past end of memory in rule printing (Peter Eisentraut)

- Fix array slicing of int2vector and oidvector values (Tom Lane)

  Expressions of this kind are now implicitly promoted to regular int2 or oid arrays.

- Fix incorrect behaviors when using a SQL-standard, simple GMT offset timezone (Tom Lane)

  In some cases, the system would use the simple GMT offset value when it should have used the regular timezone setting that had prevailed before the simple offset was selected. This change also causes thetimeofday function to honor the simple GMT offset zone.

- Prevent possible misbehavior when logging translations of Windows error codes (Tom Lane)

- Properly quote generated command lines in pg_ctl (Naoya Anzai and Tom Lane)

  This fix applies only to Windows.

- Fix pg_dumpall to work when a source database sets default_transaction_read_only via ALTER DATABASE SET (Kevin Grittner)

  Previously, the generated script would fail during restore.

- Make ecpg search for quoted cursor names case-sensitively (Zoltán Böszörményi)

- Fix ecpg's processing of lists of variables declared varchar (Zoltán Böszörményi)

- Make contrib/lo defend against incorrect trigger definitions (Marc Cousin)

- Update time zone data files to tzdata release 2013h for DST law changes in Argentina, Brazil, Jordan, Libya, Liechtenstein, Morocco, and Palestine. Also, new timezone abbreviations WIB, WIT, WITA for Indonesia.

---

**E.124. Release 9.1.10**

**Release date:** 2013-10-10

This release contains a variety of fixes from 9.1.9. For information about new features in the 9.1 major release, see Section E.134.
E.124.1. Migration to Version 9.1.10

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.9, see Section E.125.

E.124.2. Changes

- Prevent corruption of multi-byte characters when attempting to case-fold identifiers (Andrew Dunstan)

  PostgreSQL case-folds non-ASCII characters only when using a single-byte server encoding.

- Fix checkpoint memory leak in background writer when wal_level = hot_standby (Naoya Anzai)

- Fix memory leak caused by lo_open() failure (Heikki Linnakangas)

- Fix memory overcommit bug when work_mem is using more than 24GB of memory (Stephen Frost)

- Serializable snapshot fixes (Kevin Grittner, Heikki Linnakangas)

- Fix deadlock bug in libpq when using SSL (Stephen Frost)

- Fix possible SSL state corruption in threaded libpq applications (Nick Phillips, Stephen Frost)

- Properly compute row estimates for boolean columns containing many NULL values (Andrew Gierth)

  Previously tests like col IS NOT TRUE and col IS NOT FALSE did not properly factor in NULL values when estimating plan costs.

- Prevent pushing down WHERE clauses into unsafe UNION/INTERSECT subqueries (Tom Lane)

  Subqueries of a UNION or INTERSECT that contain set-returning functions or volatile functions in their SELECT lists could be improperly optimized, leading to run-time errors or incorrect query results.

- Fix rare case of “failed to locate grouping columns” planner failure (Tom Lane)

- Fix pg_dump of foreign tables with dropped columns (Andrew Dunstan)

  Previously such cases could cause a pg_upgrade error.

- Reorder pg_dump processing of extension-related rules and event triggers (Joe Conway)

- Force dumping of extension tables if specified by pg_dump -t or -n (Joe Conway)

- Improve view dumping code's handling of dropped columns in referenced tables (Tom Lane)

- Fix pg_restore -l with the directory archive to display the correct format name (Fujii Masao)

- Properly record index comments created using UNIQUE and PRIMARY KEY syntax (Andres Freund)

  This fixes a parallel pg_restore failure.

- Properly guarantee transmission of WAL files before clean switchover (Fujii Masao)

  Previously, the streaming replication connection might close before all WAL files had been replayed on the standby.
• Fix WAL segment timeline handling during recovery (Mitsumasa Kondo, Heikki Linnakangas)

  WAL file recycling during standby recovery could lead to premature recovery completion, resulting in data loss.

• Fix REINDEX TABLE and REINDEX DATABASE to properly revalidate constraints and mark invalidated indexes as valid (Noah Misch)

  REINDEX INDEX has always worked properly.

• Fix possible deadlock during concurrent CREATE INDEX CONCURRENTLY operations (Tom Lane)

• Fix regexp_matches() handling of zero-length matches (Jeevan Chalke)

  Previously, zero-length matches like '^' could return too many matches.

• Fix crash for overly-complex regular expressions (Heikki Linnakangas)

• Fix regular expression match failures for back references combined with non-greedy quantifiers (Jeevan Chalke)

• Prevent CREATE FUNCTION from checking SET variables unless function body checking is enabled (Tom Lane)

• Allow ALTER DEFAULT PRIVILEGES to operate on schemas without requiring CREATE permission (Tom Lane)

• Loosen restriction on keywords used in queries (Tom Lane)

  Specifically, lessen keyword restrictions for role names, language names, EXPLAIN and COPY options, and SET values. This allows COPY ... (FORMAT BINARY) to work as expected; previously BINARY needed to be quoted.

• Fix pgp_pub_decrypt() so it works for secret keys with passwords (Marko Kreen)

• Make pg_upgrade use pg_dump --quote-all-identifiers to avoid problems with keyword changes between releases (Tom Lane)

• Remove rare inaccurate warning during vacuum of index-less tables (Heikki Linnakangas)

• Ensure that VACUUM ANALYZE still runs the ANALYZE phase if its attempt to truncate the file is cancelled due to lock conflicts (Kevin Grittner)

• Avoid possible failure when performing transaction control commands (e.g ROLLBACK) in prepared queries (Tom Lane)

• Ensure that floating-point data input accepts standard spellings of “infinity” on all platforms (Tom Lane)

  The C99 standard says that allowable spellings are inf, +inf, -inf, infinity, +infinity, and -infinity. Make sure we recognize these even if the platform's strtod function doesn't.

• Expand ability to compare rows to records and arrays (Rafal Rzpecki, Tom Lane)

• Update time zone data files to tzdata release 2013d for DST law changes in Israel, Morocco, Palestine, and Paraguay. Also, historical zone data corrections for Macquarie Island.

E.125. Release 9.1.9

Release date: 2013-04-04
This release contains a variety of fixes from 9.1.8. For information about new features in the 9.1 major release, see Section E.134.

E.125.1. Migration to Version 9.1.9

A dump/restore is not required for those running 9.1.X.

However, this release corrects several errors in management of GiST indexes. After installing this update, it is advisable to REINDEX any GiST indexes that meet one or more of the conditions described below.

Also, if you are upgrading from a version earlier than 9.1.6, see Section E.128.

E.125.2. Changes

• Fix insecure parsing of server command-line switches (Mitsumasa Kondo, Kyotaro Horiguchi)

A connection request containing a database name that begins with “-” could be crafted to damage or destroy files within the server's data directory, even if the request is eventually rejected. (CVE-2013-1899)

• Reset OpenSSL randomness state in each postmaster child process (Marko Kreen)

This avoids a scenario wherein random numbers generated by contrib/pgcrypto functions might be relatively easy for another database user to guess. The risk is only significant when the postmaster is configured with ssl = on but most connections don't use SSL encryption. (CVE-2013-1900)

• Make REPLICATION privilege checks test current user not authenticated user (Noah Misch)

An unprivileged database user could exploit this mistake to call pg_start_backup() or pg_stop_backup(), thus possibly interfering with creation of routine backups. (CVE-2013-1901)

• Fix GiST indexes to not use “fuzzy” geometric comparisons when it's not appropriate to do so (Alexander Korotkov)

The core geometric types perform comparisons using “fuzzy” equality, but gist_box_same must do exact comparisons, else GiST indexes using it might become inconsistent. After installing this update, users should REINDEX any GiST indexes on box, polygon, circle, or point columns, since all of these use gist_box_same.

• Fix erroneous range-union and penalty logic in GiST indexes that use contrib/btree_gist for variable-width data types, that is text, bytea, bit, and numeric columns (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in useless index bloat. Users are advised to REINDEX such indexes after installing this update.

• Fix bugs in GiST page splitting code for multi-column indexes (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in indexes that are unnecessarily inefficient to search. Users are advised to REINDEX multi-column GiST indexes after installing this update.

• Fix gist_point_consistent to handle fuzziness consistently (Alexander Korotkov)

Index scans on GiST indexes on point columns would sometimes yield results different from a sequential scan, because gist_point_consistent disagreed with the underlying operator code about whether to do comparisons exactly or fuzzily.
• Fix buffer leak in WAL replay (Heikki Linnakangas)

This bug could result in “incorrect local pin count” errors during replay, making recovery impossible.

• Fix race condition in DELETE RETURNING (Tom Lane)

Under the right circumstances, DELETE RETURNING could attempt to fetch data from a shared buffer that the current process no longer has any pin on. If some other process changed the buffer meanwhile, this would lead to garbage RETURNING output, or even a crash.

• Fix infinite-loop risk in regular expression compilation (Tom Lane, Don Porter)

• Fix potential null-pointer dereference in regular expression compilation (Tom Lane)

• Fix to_char() to use ASCII-only case-folding rules where appropriate (Tom Lane)

This fixes misbehavior of some template patterns that should be locale-independent, but mishandled “I” and “i” in Turkish locales.

• Fix unwanted rejection of timestamp 1999-12-31 24:00:00 (Tom Lane)

• Fix logic error when a single transaction does UNLISTEN then LISTEN (Tom Lane)

The session wound up not listening for notify events at all, though it surely should listen in this case.

• Fix possible planner crash after columns have been added to a view that’s depended on by another view (Tom Lane)

• Remove useless “picksplit doesn’t support secondary split” log messages (Josh Hansen, Tom Lane)

This message seems to have been added in expectation of code that was never written, and probably never will be, since GiST’s default handling of secondary splits is actually pretty good. So stop nagging end users about it.

• Fix possible failure to send a session’s last few transaction commit/abort counts to the statistics collector (Tom Lane)

• Eliminate memory leaks in PL/Perl’s spi_prepare() function (Alex Hunsaker, Tom Lane)

• Fix pg_dumpall to handle database names containing “=” correctly (Heikki Linnakangas)

• Ignore invalid indexes in pg_dump and pg_upgrade (Michael Paquier, Bruce Momjian)

Dumping invalid indexes can cause problems at restore time, for example if the reason the index creation failed was because it tried to enforce a uniqueness condition not satisfied by the table’s data. Also, if the index creation is in fact still in progress, it seems reasonable to consider it to be an uncommitted DDL change, which pg_dump wouldn’t be expected to dump anyway. pg_upgrade now also skips invalid indexes rather than failing.

• In pg_basebackup, include only the current server version’s subdirectory when backing up a table-space (Heikki Linnakangas)

• Add a server version check in pg_basebackup and pg_receivexlog, so they fail cleanly with version combinations that won’t work (Heikki Linnakangas)

• Fix contrib/pg_trgm’s similarity() function to return zero for trigram-less strings (Tom Lane)

Previously it returned NaN due to internal division by zero.
• Update time zone data files to tzdata release 2013b for DST law changes in Chile, Haiti, Morocco, Paraguay, and some Russian areas. Also, historical zone data corrections for numerous places.

Also, update the time zone abbreviation files for recent changes in Russia and elsewhere: CHOT, GET, IRKT, KGT, KRAT, MAGT, MAWT, MSK, NOVT, OMST, TKT, VLAT, WST, YAKT, YEKT now follow their current meanings, and VOLT (Europe/Volgograd) and MIST (Antarctica/Macquarie) are added to the default abbreviations list.

E.126. Release 9.1.8

Release date: 2013-02-07

This release contains a variety of fixes from 9.1.7. For information about new features in the 9.1 major release, see Section E.134.

E.126.1. Migration to Version 9.1.8

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.6, see Section E.128.

E.126.2. Changes

• Prevent execution of enum_recv from SQL (Tom Lane)

The function was misdeclared, allowing a simple SQL command to crash the server. In principle an attacker might be able to use it to examine the contents of server memory. Our thanks to Sumit Soni (via Secunia SVCRP) for reporting this issue. (CVE-2013-0255)

• Fix multiple problems in detection of when a consistent database state has been reached during WAL replay (Fujii Masao, Heikki Linnakangas, Simon Riggs, Andres Freund)

• Update minimum recovery point when truncating a relation file (Heikki Linnakangas)

Once data has been discarded, it’s no longer safe to stop recovery at an earlier point in the timeline.

• Fix recycling of WAL segments after changing recovery target timeline (Heikki Linnakangas)

• Fix missing cancellations in hot standby mode (Noah Misch, Simon Riggs)

The need to cancel conflicting hot-standby queries would sometimes be missed, allowing those queries to see inconsistent data.

• Prevent recovery pause feature from pausing before users can connect (Tom Lane)

• Fix SQL grammar to allow subscripting or field selection from a sub-SELECT result (Tom Lane)

• Fix performance problems with autovacuum truncation in busy workloads (Jan Wieck)

Truncation of empty pages at the end of a table requires exclusive lock, but autovacuum was coded to fail (and release the table lock) when there are conflicting lock requests. Under load, it is easily possible that truncation would never occur, resulting in table bloat. Fix by performing a partial truncation, releasing the lock, then attempting to re-acquire the lock and continue. This fix also greatly reduces the average time before autovacuum releases the lock after a conflicting request arrives.

• Prevent race conditions when scanning pg_tablespace (Stephen Frost, Tom Lane)

CREATE DATABASE and DROP DATABASE could misbehave if there were concurrent updates of pg_tablespace entries.
- Prevent `DROP OWNED` from trying to drop whole databases or tablespaces (Álvaro Herrera)
  For safety, ownership of these objects must be reassigned, not dropped.
- Fix error in `vacuum_freeze_table_age` implementation (Andres Freund)
  In installations that have existed for more than `vacuum_freeze_min_age` transactions, this mistake prevented autovacuum from using partial-table scans, so that a full-table scan would always happen instead.
- Prevent misbehavior when a `RowExpr` or `XmlExpr` is parse-analyzed twice (Andres Freund, Tom Lane)
  This mistake could be user-visible in contexts such as `CREATE TABLE LIKE INCLUDING INDEXES`.
- Improve defenses against integer overflow in hashtable sizing calculations (Jeff Davis)
- Fix failure to ignore leftover temporary tables after a server crash (Tom Lane)
- Reject out-of-range dates in `to_date()` (Hitoshi Harada)
- Fix `pg_extension_config_dump()` to handle extension-update cases properly (Tom Lane)
  This function will now replace any existing entry for the target table, making it usable in extension update scripts.
- Fix PL/Python's handling of functions used as triggers on multiple tables (Andres Freund)
- Ensure that non-ASCII prompt strings are translated to the correct code page on Windows (Alexander Law, Noah Misch)
  This bug affected psql and some other client programs.
- Fix possible crash in psql's `?` command when not connected to a database (Meng Qingzhong)
- Fix possible error if a relation file is removed while `pg_basebackup` is running (Heikki Linnakangas)
- Make `pg_dump` exclude data of unlogged tables when running on a hot-standby server (Magnus Hagander)
  This would fail anyway because the data is not available on the standby server, so it seems most convenient to assume `--no-unlogged-table-data` automatically.
- Fix `pg_upgrade` to deal with invalid indexes safely (Bruce Momjian)
- Fix one-byte buffer overrun in `libpq`'s `PQprintTuples` (Xi Wang)
  This ancient function is not used anywhere by PostgreSQL itself, but it might still be used by some client code.
- Make ecpglib use translated messages properly (Chen Huajun)
- Properly install `ecpg_compat` and `pgtypes` libraries on MSVC (Jiang Guiqing)
- Include our version of `isinf()` in libecpg if it's not provided by the system (Jiang Guiqing)
- Rearrange configure's tests for supplied functions so it is not fooled by bogus exports from libedit/libreadline (Christoph Berg)
- Ensure Windows build number increases over time (Magnus Hagander)
- Make `pgxs` build executables with the right `.exe` suffix when cross-compiling for Windows (Zoltan Boszormenyi)
• Add new timezone abbreviation FET (Tom Lane)

This is now used in some eastern-European time zones.

E.127. Release 9.1.7

Release date: 2012-12-06

This release contains a variety of fixes from 9.1.6. For information about new features in the 9.1 major release, see Section E.134.

E.127.1. Migration to Version 9.1.7

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.6, see Section E.128.

E.127.2. Changes

• Fix multiple bugs associated with CREATE INDEX CONCURRENTLY (Andres Freund, Tom Lane)

  Fix CREATE INDEX CONCURRENTLY to use in-place updates when changing the state of an index’s pg_index row. This prevents race conditions that could cause concurrent sessions to miss updating the target index, thus resulting in corrupt concurrently-created indexes.

  Also, fix various other operations to ensure that they ignore invalid indexes resulting from a failed CREATE INDEX CONCURRENTLY command. The most important of these is VACUUM, because an auto-vacuum could easily be launched on the table before corrective action can be taken to fix or remove the invalid index.

• Fix buffer locking during WAL replay (Tom Lane)

  The WAL replay code was insufficiently careful about locking buffers when replaying WAL records that affect more than one page. This could result in hot standby queries transiently seeing inconsistent states, resulting in wrong answers or unexpected failures.

• Fix an error in WAL generation logic for GIN indexes (Tom Lane)

  This could result in index corruption, if a torn-page failure occurred.

• Properly remove startup process’s virtual XID lock when promoting a hot standby server to normal running (Simon Riggs)

  This oversight could prevent subsequent execution of certain operations such as CREATE INDEX CONCURRENTLY.

• Avoid bogus “out-of-sequence timeline ID” errors in standby mode (Heikki Linnakangas)

• Prevent the postmaster from launching new child processes after it’s received a shutdown signal (Tom Lane)

  This mistake could result in shutdown taking longer than it should, or even never completing at all without additional user action.

• Avoid corruption of internal hash tables when out of memory (Hitoshi Harada)

• Prevent file descriptors for dropped tables from being held open past transaction end (Tom Lane)

  This should reduce problems with long-since-dropped tables continuing to occupy disk space.
• Prevent database-wide crash and restart when a new child process is unable to create a pipe for its latch (Tom Lane)

Although the new process must fail, there is no good reason to force a database-wide restart, so avoid that. This improves robustness when the kernel is nearly out of file descriptors.

• Fix planning of non-strict equivalence clauses above outer joins (Tom Lane)

The planner could derive incorrect constraints from a clause equating a non-strict construct to something else, for example \texttt{WHERE COALESCE(foo, 0) = 0} when \texttt{foo} is coming from the nullable side of an outer join.

• Fix \texttt{SELECT DISTINCT} with index-optimized \texttt{MIN/MAX} on an inheritance tree (Tom Lane)

The planner would fail with “failed to re-find MinMaxAggInfo record” given this combination of factors.

• Improve planner’s ability to prove exclusion constraints from equivalence classes (Tom Lane)

• Fix partial-row matching in hashed subplans to handle cross-type cases correctly (Tom Lane)

This affects multicol \texttt{NOT IN} subplans, such as \texttt{WHERE (a, b) NOT IN (SELECT x, y FROM ...)} when for instance \texttt{b} and \texttt{y} are \texttt{int4} and \texttt{int8} respectively. This mistake led to wrong answers or crashes depending on the specific datatypes involved.

• Acquire buffer lock when re-fetching the old tuple for an \texttt{AFTER ROW UPDATE/DELETE} trigger (Andres Freund)

In very unusual circumstances, this oversight could result in passing incorrect data to a trigger \texttt{WHEN} condition, or to the precheck logic for a foreign-key enforcement trigger. That could result in a crash, or in an incorrect decision about whether to fire the trigger.

• Fix \texttt{ALTER COLUMN TYPE} to handle inherited check constraints properly (Pavan Deolasee)

This worked correctly in pre-8.4 releases, and now works correctly in 8.4 and later.

• Fix \texttt{ALTER EXTENSION SET SCHEMA}’s failure to move some subsidiary objects into the new schema (Álvaro Herrera, Dimitri Fontaine)

• Fix \texttt{REASSIGN OWNED} to handle grants on tablespaces (Álvaro Herrera)

• Ignore incorrect \texttt{pg_attribute} entries for system columns for views (Tom Lane)

Views do not have any system columns. However, we forgot to remove such entries when converting a table to a view. That’s fixed properly for 9.3 and later, but in previous branches we need to defend against existing mis-converted views.

• Fix rule printing to dump \texttt{INSERT INTO table DEFAULT VALUES} correctly (Tom Lane)

• Guard against stack overflow when there are too many \texttt{UNION/INTERSECT/EXCEPT} clauses in a query (Tom Lane)

• Prevent platform-dependent failures when dividing the minimum possible integer value by -1 (Xi Wang, Tom Lane)

• Fix possible access past end of string in date parsing (Hitoshi Harada)

• Fix failure to advance XID epoch if XID wraparound happens during a checkpoint and \texttt{wal_level} is \texttt{hot_standby} (Tom Lane, Andres Freund)

While this mistake had no particular impact on PostgreSQL itself, it was bad for applications that rely on \texttt{txid_current()} and related functions: the TXID value would appear to go backwards.
• Fix display of pg_stat_replication.sync_state at a page boundary (Kyotaro Horiguchi)

• Produce an understandable error message if the length of the path name for a Unix-domain socket exceeds the platform-specific limit (Tom Lane, Andrew Dunstan)

Formerly, this would result in something quite unhelpful, such as “Non-recoverable failure in name resolution”.

• Fix memory leaks when sending composite column values to the client (Tom Lane)

• Make pg_ctl more robust about reading the postmaster.pid file (Heikki Linnakangas)

Fix race conditions and possible file descriptor leakage.

• Fix possible crash in psql if incorrectly-encoded data is presented and the client_encoding setting is a client-only encoding, such as SJIS (Jiang Guiqing)

• Make pg_dump dump SEQUENCE SET items in the data not pre-data section of the archive (Tom Lane)

This change fixes dumping of sequences that are marked as extension configuration tables.

• Fix bugs in the restore.sql script emitted by pg_dump in tar output format (Tom Lane)

The script would fail outright on tables whose names include upper-case characters. Also, make the script capable of restoring data in --inserts mode as well as the regular COPY mode.

• Fix pg_restore to accept POSIX-conformant tar files (Brian Weaver, Tom Lane)

The original coding of pg_dump's tar output mode produced files that are not fully conformant with the POSIX standard. This has been corrected for version 9.3. This patch updates previous branches so that they will accept both the incorrect and the corrected formats, in hopes of avoiding compatibility problems when 9.3 comes out.

• Fix tar files emitted by pg_basebackup to be POSIX conformant (Brian Weaver, Tom Lane)

• Fix pg_resetxlog to locate postmaster.pid correctly when given a relative path to the data directory (Tom Lane)

This mistake could lead to pg_resetxlog not noticing that there is an active postmaster using the data directory.

• Fix libpq's lo_import() and lo_export() functions to report file I/O errors properly (Tom Lane)

• Fix ecpg's processing of nested structure pointer variables (Muhammad Usama)

• Fix ecpg's ecpg_get_data function to handle arrays properly (Michael Meskes)

• Make contrib/pageinspect's btree page inspection functions take buffer locks while examining pages (Tom Lane)

• Ensure that make install for an extension creates the extension installation directory (Cédrick Villemain)

Previously, this step was missed if MODULEDIR was set in the extension's Makefile.

• Fix pgxs support for building loadable modules on AIX (Tom Lane)

Building modules outside the original source tree didn't work on AIX.

• Update time zone data files to tzdata release 2012j for DST law changes in Cuba, Israel, Jordan, Libya, Palestine, Western Samoa, and portions of Brazil.
E.128. Release 9.1.6

Release date: 2012-09-24

This release contains a variety of fixes from 9.1.5. For information about new features in the 9.1 major release, see Section E.134.

E.128.1. Migration to Version 9.1.6

A dump/restore is not required for those running 9.1.X.

However, you may need to perform REINDEX operations to recover from the effects of the data corruption bug described in the first changelog item below.

Also, if you are upgrading from a version earlier than 9.1.4, see Section E.130.

E.128.2. Changes

• Fix persistence marking of shared buffers during WAL replay (Jeff Davis)

This mistake can result in buffers not being written out during checkpoints, resulting in data corruption if the server later crashes without ever having written those buffers. Corruption can occur on any server following crash recovery, but it is significantly more likely to occur on standby slave servers since those perform much more WAL replay. There is a low probability of corruption of btree and GIN indexes. There is a much higher probability of corruption of table “visibility maps”. Fortunately, visibility maps are non-critical data in 9.1, so the worst consequence of such corruption in 9.1 installations is transient inefficiency of vacuuming. Table data proper cannot be corrupted by this bug.

While no index corruption due to this bug is known to have occurred in the field, as a precautionary measure it is recommended that production installations REINDEX all btree and GIN indexes at a convenient time after upgrading to 9.1.6.

Also, if you intend to do an in-place upgrade to 9.2.X, before doing so it is recommended to perform a VACUUM of all tables while having vacuum_freeze_table_age set to zero. This will ensure that any lingering wrong data in the visibility maps is corrected before 9.2.X can depend on it. vacuum_cost_delay can be adjusted to reduce the performance impact of vacuuming, while causing it to take longer to finish.

• Fix planner’s assignment of executor parameters, and fix executor’s rescan logic for CTE plan nodes (Tom Lane)

These errors could result in wrong answers from queries that scan the same WITH subquery multiple times.

• Fix misbehavior when default_transaction_isolation is set to serializable (Kevin Grittner, Tom Lane, Heikki Linnakangas)

Symptoms include crashes at process start on Windows, and crashes in hot standby operation.

• Improve selectivity estimation for text search queries involving prefixes, i.e. word:* patterns (Tom Lane)

• Improve page-splitting decisions in GiST indexes (Alexander Korotkov, Robert Haas, Tom Lane)

Multi-column GiST indexes might suffer unexpected bloat due to this error.

• Fix cascading privilege revoke to stop if privileges are still held (Tom Lane)
If we revoke a grant option from some role \( X \), but \( X \) still holds that option via a grant from someone else, we should not recursively revoke the corresponding privilege from role(s) \( Y \) that \( X \) had granted it to.

- Disallow extensions from containing the schema they are assigned to (Thom Brown)
  This situation creates circular dependencies that confuse pg_dump and probably other things. It’s confusing for humans too, so disallow it.

- Improve error messages for Hot Standby misconfiguration errors (Gurjeet Singh)

- Make configure probe for `mbstowcs_l` (Tom Lane)
  This fixes build failures on some versions of AIX.

- Fix handling of `SIGFPE` when PL/Perl is in use (Andres Freund)
  Perl resets the process's `SIGFPE` handler to `SIG_IGN`, which could result in crashes later on. Restore the normal Postgres signal handler after initializing PL/Perl.

- Prevent PL/Perl from crashing if a recursive PL/Perl function is redefined while being executed (Tom Lane)

- Work around possible misoptimization in PL/Perl (Tom Lane)
  Some Linux distributions contain an incorrect version of `pthread.h` that results in incorrect compiled code in PL/Perl, leading to crashes if a PL/Perl function calls another one that throws an error.

- Fix bugs in contrib/pg_trgm's LIKE pattern analysis code (Fujii Masao)
  LIKE queries using a trigram index could produce wrong results if the pattern contained LIKE escape characters.

- Fix `pg_upgrade`'s handling of line endings on Windows (Andrew Dunstan)
  Previously, `pg_upgrade` might add or remove carriage returns in places such as function bodies.

- On Windows, make `pg_upgrade` use backslash path separators in the scripts it emits (Andrew Dunstan)

- Remove unnecessary dependency on `pg_config` from `pg_upgrade` (Peter Eisentraut)

- Update time zone data files to tzdata release 2012f for DST law changes in Fiji

### E.129. Release 9.1.5

**Release date:** 2012-08-17

This release contains a variety of fixes from 9.1.4. For information about new features in the 9.1 major release, see Section E.134.

### E.129.1. Migration to Version 9.1.5

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.4, see Section E.130.

### E.129.2. Changes

- Prevent access to external files/URLs via XML entity references (Noah Misch, Tom Lane)
xml_parse() would attempt to fetch external files or URLs as needed to resolve DTD and entity references in an XML value, thus allowing unprivileged database users to attempt to fetch data with the privileges of the database server. While the external data wouldn’t get returned directly to the user, portions of it could be exposed in error messages if the data didn’t parse as valid XML; and in any case the mere ability to check existence of a file might be useful to an attacker. (CVE-2012-3489)

• Prevent access to external files/URLs via contrib/xml2’s xslt_process() (Peter Eisentraut)

libxslt offers the ability to read and write both files and URLs through stylesheet commands, thus allowing unprivileged database users to both read and write data with the privileges of the database server. Disable that through proper use of libxslt’s security options. (CVE-2012-3488)

Also, remove xslt_process()’s ability to fetch documents and stylesheets from external files/URLs. While this was a documented “feature”, it was long regarded as a bad idea. The fix for CVE-2012-3489 broke that capability, and rather than expend effort on trying to fix it, we’re just going to summarily remove it.

• Prevent too-early recycling of btree index pages (Noah Misch)

When we allowed read-only transactions to skip assigning XIDs, we introduced the possibility that a deleted btree page could be recycled while a read-only transaction was still in flight to it. This would result in incorrect index search results. The probability of such an error occurring in the field seems very low because of the timing requirements, but nonetheless it should be fixed.

• Fix crash-safety bug with newly-created-or-reset sequences (Tom Lane)

If ALTER SEQUENCE was executed on a freshly created or reset sequence, and then precisely one nextval() call was made on it, and then the server crashed, WAL replay would restore the sequence to a state in which it appeared that no nextval() had been done, thus allowing the first sequence value to be returned again by the next nextval() call. In particular this could manifest for serial columns, since creation of a serial column’s sequence includes an ALTER SEQUENCE OWNED BY step.

• Fix race condition in enum-type value comparisons (Robert Haas, Tom Lane)

Comparisons could fail when encountering an enum value added since the current query started.

• Fix txid_current() to report the correct epoch when not in hot standby (Heikki Linnakangas)

This fixes a regression introduced in the previous minor release.

• Prevent selection of unsuitable replication connections as the synchronous standby (Fujii Masao)

The master might improperly choose pseudo-servers such as pg_receivexlog or pg_basebackup as the synchronous standby, and then wait indefinitely for them.

• Fix bug in startup of Hot Standby when a master transaction has many subtransactions (Andres Freund)

This mistake led to failures reported as “out-of-order XID insertion in KnownAssignedXids”.

• Ensure the backup_label file is fsync’d after pg_start_backup() (Dave Kerr)

• Fix timeout handling in walsender processes (Tom Lane)

WAL sender background processes neglected to establish a SIGALRM handler, meaning they would wait forever in some corner cases where a timeout ought to happen.

• Wake walsenders after each background flush by walwriter (Andres Freund, Simon Riggs)
This greatly reduces replication delay when the workload contains only asynchronously-committed transactions.

- **Fix LISTEN/NOTIFY to cope better with I/O problems, such as out of disk space (Tom Lane)**

  After a write failure, all subsequent attempts to send more NOTIFY messages would fail with messages like “Could not read from file "pg_notify/nnnn" at offset nnnnn: Success”.

- **Only allow autovacuum to be auto-canceled by a directly blocked process (Tom Lane)**

  The original coding could allow inconsistent behavior in some cases; in particular, an autovacuum could get canceled after less than deadlock_timeout grace period.

- **Improve logging of autovacuum cancels (Robert Haas)**

- **Fix log collector so that log_truncate_on_rotation works during the very first log rotation after server start (Tom Lane)**

- **Fix WITH attached to a nested set operation (UNION/INTERSECT/EXCEPT) (Tom Lane)**

- **Ensure that a whole-row reference to a subquery doesn’t include any extra GROUP BY or ORDER BY columns (Tom Lane)**

- **Fix dependencies generated during ALTER TABLE ... ADD CONSTRAINT USING INDEX (Tom Lane)**

  This command left behind a redundant pg_depend entry for the index, which could confuse later operations, notably ALTER TABLE ... ALTER COLUMN TYPE on one of the indexed columns.

- **Fix REASSIGN OWNED to work on extensions (Alvaro Herrera)**

- **Disallow copying whole-row references in CHECK constraints and index definitions during CREATE TABLE (Tom Lane)**

  This situation can arise in CREATE TABLE with LIKE or INHERITS. The copied whole-row variable was incorrectly labeled with the row type of the original table not the new one. Rejecting the case seems reasonable for LIKE, since the row types might well diverge later. For INHERITS we should ideally allow it, with an implicit coercion to the parent table’s row type; but that will require more work than seems safe to back-patch.

- **Fix memory leak in ARRAY(SELECT ...) subqueries (Heikki Linnakangas, Tom Lane)**

- **Fix planner to pass correct collation to operator selectivity estimators (Tom Lane)**

  This was not previously required by any core selectivity estimation function, but third-party code might need it.

- **Fix extraction of common prefixes from regular expressions (Tom Lane)**

  The code could get confused by quantified parenthesized subexpressions, such as ^(foo)?bar. This would lead to incorrect index optimization of searches for such patterns.

- **Fix bugs with parsing signed hh:mm and hh:mm:ss fields in interval constants (Amit Kapila, Tom Lane)**

- **Fix pg_dump to better handle views containing partial GROUP BY lists (Tom Lane)**

  A view that lists only a primary key column in GROUP BY, but uses other table columns as if they were grouped, gets marked as depending on the primary key. Improper handling of such primary key dependencies in pg_dump resulted in poorly-ordered dumps, which at best would be inefficient to restore and at worst could result in outright failure of a parallel pg_restore run.
• In PL/Perl, avoid setting UTF8 flag when in SQL_ASCII encoding (Alex Hunsaker, Kyotaro Horiguchi, Alvaro Herrera)

• Use Postgres' encoding conversion functions, not Python's, when converting a Python Unicode string to the server encoding in PL/Python (Jan Urbanski)

This avoids some corner-case problems, notably that Python doesn't support all the encodings Postgres does. A notable functional change is that if the server encoding is SQL_ASCII, you will get the UTF-8 representation of the string; formerly, any non-ASCII characters in the string would result in an error.

• Fix mapping of PostgreSQL encodings to Python encodings in PL/Python (Jan Urbanski)

• Report errors properly in contrib/xml2\'s xslt_process() (Tom Lane)

• Update time zone data files to tzdata release 2012e for DST law changes in Morocco and Tokelau

E.130. Release 9.1.4

Release date: 2012-06-04

This release contains a variety of fixes from 9.1.3. For information about new features in the 9.1 major release, see Section E.134.

E.130.1. Migration to Version 9.1.4

A dump/restore is not required for those running 9.1.X.

However, if you use the citext data type, and you upgraded from a previous major release by running pg_upgrade, you should run CREATE EXTENSION citext FROM unpackaged to avoid collation-related failures in citext operations. The same is necessary if you restore a dump from a pre-9.1 database that contains an instance of the citext data type. If you've already run the CREATE EXTENSION command before upgrading to 9.1.4, you will instead need to do manual catalog updates as explained in the third changelog item below.

Also, if you are upgrading from a version earlier than 9.1.2, see Section E.132.

E.130.2. Changes

• Fix incorrect password transformation in contrib/pgcrypto\'s DES crypt() function (Solar Designer)

If a password string contained the byte value 0x80, the remainder of the password was ignored, causing the password to be much weaker than it appeared. With this fix, the rest of the string is properly included in the DES hash. Any stored password values that are affected by this bug will thus no longer match, so the stored values may need to be updated. (CVE-2012-2143)

• Ignore SECURITY DEFINER and SET attributes for a procedural language\'s call handler (Tom Lane)

Applying such attributes to a call handler could crash the server. (CVE-2012-2655)

• Make contrib/citext\'s upgrade script fix collations of citext arrays and domains over citext (Tom Lane)

Release 9.1.2 provided a fix for collations of citext columns and indexes in databases upgraded or reloaded from pre-9.1 installations, but that fix was incomplete: it neglected to handle arrays and domains over citext. This release extends the module\'s upgrade script to handle these cases. As before, if you have already run the upgrade script, you\'ll need to run the collation update commands by hand instead. See the 9.1.2 release notes for more information about doing this.
• Allow numeric timezone offsets in `timestamp` input to be up to 16 hours away from UTC (Tom Lane)

Some historical time zones have offsets larger than 15 hours, the previous limit. This could result in dumped data values being rejected during reload.

• Fix timestamp conversion to cope when the given time is exactly the last DST transition time for the current timezone (Tom Lane)

This oversight has been there a long time, but was not noticed previously because most DST-using zones are presumed to have an indefinite sequence of future DST transitions.

• Fix `text` to `name` and `char` to `name` casts to perform string truncation correctly in multibyte encodings (Karl Schnaitter)

• Fix memory copying bug in `to_tsquery()` (Heikki Linnakangas)

• Ensure `txid_current()` reports the correct epoch when executed in hot standby (Simon Riggs)

• Fix planner’s handling of outer PlaceHolderVars within subqueries (Tom Lane)

This bug concerns sub-SELECTs that reference variables coming from the nullable side of an outer join of the surrounding query. In 9.1, queries affected by this bug would fail with “ERROR: Upper-level PlaceHolderVar found where not expected”. But in 9.0 and 8.4, you’d silently get possibly-wrong answers, since the value transmitted into the subquery wouldn’t go to null when it should.

• Fix planning of `UNION ALL` subqueries with output columns that are not simple variables (Tom Lane)

Planning of such cases got noticeably worse in 9.1 as a result of a misguided fix for “MergeAppend child’s targetlist doesn’t match MergeAppend” errors. Revert that fix and do it another way.

• Fix slow session startup when `pg_attribute` is very large (Tom Lane)

If `pg_attribute` exceeds one-fourth of `shared_buffers`, cache rebuilding code that is sometimes needed during session start would trigger the synchronized-scan logic, causing it to take many times longer than normal. The problem was particularly acute if many new sessions were starting at once.

• Ensure sequential scans check for query cancel reasonably often (Merlin Moncure)

A scan encountering many consecutive pages that contain no live tuples would not respond to interrupts meanwhile.

• Ensure the Windows implementation of `PGSemaphoreLock()` clears `ImmediateInterruptOK` before returning (Tom Lane)

This oversight meant that a query-cancel interrupt received later in the same query could be accepted at an unsafe time, with unpredictable but not good consequences.

• Show whole-row variables safely when printing views or rules (Abbas Butt, Tom Lane)

Corner cases involving ambiguous names (that is, the name could be either a table or column name of the query) were printed in an ambiguous way, risking that the view or rule would be interpreted differently after dump and reload. Avoid the ambiguous case by attaching a no-op cast.

• Fix `COPY FROM` to properly handle null marker strings that correspond to invalid encoding (Tom Lane)

A null marker string such as `E'\0'` should work, and did work in the past, but the case got broken in 8.4.
• Fix EXPLAIN VERBOSE for writable CTEs containing RETURNING clauses (Tom Lane)

• Fix PREPARE TRANSACTION to work correctly in the presence of advisory locks (Tom Lane)

  Historically, PREPARE TRANSACTION has simply ignored any session-level advisory locks the session holds, but this case was accidentally broken in 9.1.

• Fix truncation of unlogged tables (Robert Haas)

• Ignore missing schemas during non-interactive assignments of search_path (Tom Lane)

  This re-aligns 9.1's behavior with that of older branches. Previously 9.1 would throw an error for nonexistent schemas mentioned in search_path settings obtained from places such as ALTER DATABASE SET.

• Fix bugs with temporary or transient tables used in extension scripts (Tom Lane)

  This includes cases such as a rewriting of ALTER TABLE within an extension update script, since that uses a transient table behind the scenes.

• Ensure autovacuum worker processes perform stack depth checking properly (Heikki Linnakangas)

  Previously, infinite recursion in a function invoked by auto-ANALYZE could crash worker processes.

• Fix logging collector to not lose log coherency under high load (Andrew Dunstan)

  The collector previously could fail to reassemble large messages if it got too busy.

• Fix logging collector to ensure it will restart file rotation after receiving SIGHUP (Tom Lane)

• Fix “too many LWLocks taken” failure in GiST indexes (Heikki Linnakangas)

• Fix WAL replay logic for GIN indexes to not fail if the index was subsequently dropped (Tom Lane)

• Correctly detect SSI conflicts of prepared transactions after a crash (Dan Ports)

• Avoid synchronous replication delay when committing a transaction that only modified temporary tables (Heikki Linnakangas)

  In such a case the transaction's commit record need not be flushed to standby servers, but some of the code didn't know that and waited for it to happen anyway.

• Fix error handling in pg_basebackup (Thomas Ogrisegg, Fujii Masao)

• Fix walsender to not go into a busy loop if connection is terminated (Fujii Masao)

• Fix memory leak in PL/pgSQL's RETURN NEXT command (Joe Conway)

• Fix PL/pgSQL's GET DIAGNOSTICS command when the target is the function's first variable (Tom Lane)

• Ensure that PL/Perl package-qualifies the _TD variable (Alex Hunsaker)

  This bug caused trigger invocations to fail when they are nested within a function invocation that changes the current package.

• Fix PL/Python functions returning composite types to accept a string for their result value (Jan Urbanski)

  This case was accidentally broken by the 9.1 additions to allow a composite result value to be supplied in other formats, such as dictionaries.
• Fix potential access off the end of memory in psql's expanded display (\x) mode (Peter Eisentraut)

• Fix several performance problems in pg_dump when the database contains many objects (Jeff Janes, Tom Lane)

  pg_dump could get very slow if the database contained many schemas, or if many objects are in dependency loops, or if there are many owned sequences.

• Fix memory and file descriptor leaks in pg_restore when reading a directory-format archive (Peter Eisentraut)

• Fix pg_upgrade for the case that a database stored in a non-default tablespace contains a table in the cluster's default tablespace (Bruce Momjian)

• Fix memory and file descriptor leaks in pg_restore when reading a directory-format archive (Peter Eisentraut)

• Fix memory and file descriptor leaks in pg_restore when reading a directory-format archive (Peter Eisentraut)

E.131. Release 9.1.3

Release date: 2012-02-27

This release contains a variety of fixes from 9.1.2. For information about new features in the 9.1 major release, see Section E.134.

E.131.1. Migration to Version 9.1.3

A dump/restore is not required for those running 9.1.X.

However, if you are upgrading from a version earlier than 9.1.2, see Section E.132.

E.131.2. Changes

• Require execute permission on the trigger function for CREATE TRIGGER (Robert Haas)

  This missing check could allow another user to execute a trigger function with forged input data, by installing it on a table he owns. This is only of significance for trigger functions marked SECURITY DEFINER, since otherwise trigger functions run as the table owner anyway. (CVE-2012-0866)

• Remove arbitrary limitation on length of common name in SSL certificates (Heikki Linnakangas)

  Both libpq and the server truncated the common name extracted from an SSL certificate at 32 bytes. Normally this would cause nothing worse than an unexpected verification failure, but there are some
rather-implausible scenarios in which it might allow one certificate holder to impersonate another. The victim would have to have a common name exactly 32 bytes long, and the attacker would have to persuade a trusted CA to issue a certificate in which the common name has that string as a prefix. Impersonating a server would also require some additional exploit to redirect client connections. (CVE-2012-0867)

- Convert newlines to spaces in names written in pg_dump comments (Robert Haas)

    pg_dump was incautious about sanitizing object names that are emitted within SQL comments in its output script. A name containing a newline would at least render the script syntactically incorrect. Maliciously crafted object names could present a SQL injection risk when the script is reloaded. (CVE-2012-0868)

- Fix btree index corruption from insertions concurrent with vacuuming (Tom Lane)

    An index page split caused by an insertion could sometimes cause a concurrently-running VACUUM to miss removing index entries that it should remove. After the corresponding table rows are removed, the dangling index entries would cause errors (such as “could not read block N in file ...”) or worse, silently wrong query results after unrelated rows are re-inserted at the now-free table locations. This bug has been present since release 8.2, but occurs so infrequently that it was not diagnosed until now. If you have reason to suspect that it has happened in your database, reindexing the affected index will fix things.

- Fix transient zeroing of shared buffers during WAL replay (Tom Lane)

    The replay logic would sometimes zero and refill a shared buffer, so that the contents were transiently invalid. In hot standby mode this can result in a query that's executing in parallel seeing garbage data. Various symptoms could result from that, but the most common one seems to be “invalid memory alloc request size”.

- Fix handling of data-modifying WITH subplans in READ COMMITTED rechecking (Tom Lane)

    A WITH clause containing INSERT/UPDATE/DELETE would crash if the parent UPDATE or DELETE command needed to be re-evaluated at one or more rows due to concurrent updates in READ COMMITTED mode.

- Fix corner case in SSI transaction cleanup (Dan Ports)

    When finishing up a read-write serializable transaction, a crash could occur if all remaining active serializable transactions are read-only.

- Fix postmaster to attempt restart after a hot-standby crash (Tom Lane)

    A logic error caused the postmaster to terminate, rather than attempt to restart the cluster, if any backend process crashed while operating in hot standby mode.

- Fix CLUSTER/VACUUM FULL handling of toast values owned by recently-updated rows (Tom Lane)

    This oversight could lead to “duplicate key value violates unique constraint” errors being reported against the toast table's index during one of these commands.

- Update per-column permissions, not only per-table permissions, when changing table owner (Tom Lane)

    Failure to do this meant that any previously granted column permissions were still shown as having been granted by the old owner. This meant that neither the new owner nor a superuser could revoke the now-untraceable-to-table-owner permissions.

- Support foreign data wrappers and foreign servers in REASSIGN OWNED (Alvaro Herrera)
This command failed with “unexpected classid” errors if it needed to change the ownership of any such objects.

• Allow non-existent values for some settings in `ALTER USER/DATABASE SET` (Heikki Linnakangas)

  Allow `default_text_search_config`, `default_tablespace`, and `temp_tablespaces` to be set to names that are not known. This is because they might be known in another database where the setting is intended to be used, or for the tablespace cases because the tablespace might not be created yet. The same issue was previously recognized for `search_path`, and these settings now act like that one.

• Fix “unsupported node type” error caused by `COLLATE` in an `INSERT` expression (Tom Lane)

• Avoid crashing when we have problems deleting table files post-commit (Tom Lane)

  Dropping a table should lead to deleting the underlying disk files only after the transaction commits. In event of failure then (for instance, because of wrong file permissions) the code is supposed to just emit a warning message and go on, since it’s too late to abort the transaction. This logic got broken as of release 8.4, causing such situations to result in a PANIC and an unrestartable database.

• Recover from errors occurring during WAL replay of `DROP TABLESPACE` (Tom Lane)

  Replay will attempt to remove the tablespace’s directories, but there are various reasons why this might fail (for example, incorrect ownership or permissions on those directories). Formerly the replay code would panic, rendering the database unrestartable without manual intervention. It seems better to log the problem and continue, since the only consequence of failure to remove the directories is some wasted disk space.

• Fix race condition in logging `AccessExclusiveLocks` for hot standby (Simon Riggs)

  Sometimes a lock would be logged as being held by “transaction zero”. This is at least known to produce assertion failures on slave servers, and might be the cause of more serious problems.

• Track the OID counter correctly during WAL replay, even when it wraps around (Tom Lane)

  Previously the OID counter would remain stuck at a high value until the system exited replay mode. The practical consequences of that are usually nil, but there are scenarios wherein a standby server that’s been promoted to master might take a long time to advance the OID counter to a reasonable value once values are needed.

• Prevent emitting misleading “consistent recovery state reached” log message at the beginning of crash recovery (Heikki Linnakangas)

• Fix initial value of `pg_stat_replication.replay_location` (Fujii Masao)

  Previously, the value shown would be wrong until at least one WAL record had been replayed.

• Fix regular expression back-references with * attached (Tom Lane)

  Rather than enforcing an exact string match, the code would effectively accept any string that satisfies the pattern sub-expression referenced by the back-reference symbol.

  A similar problem still afflicts back-references that are embedded in a larger quantified expression, rather than being the immediate subject of the quantifier. This will be addressed in a future PostgreSQL release.

• Fix recently-introduced memory leak in processing of `inet/cidr` values (Heikki Linnakangas)

  A patch in the December 2011 releases of PostgreSQL caused memory leakage in these operations, which could be significant in scenarios such as building a btree index on such a column.
• Fix planner's ability to push down index-expression restrictions through \texttt{UNION ALL} (Tom Lane)
  This type of optimization was inadvertently disabled by a fix for another problem in 9.1.2.
• Fix planning of \texttt{WITH} clauses referenced in \texttt{UPDATE/DELETE} on an inherited table (Tom Lane)
  This bug led to "could not find plan for CTE" failures.
• Fix GIN cost estimation to handle \texttt{column IN (...) index conditions} (Marti Raudsepp)
  This oversight would usually lead to crashes if such a condition could be used with a GIN index.
• Prevent assertion failure when exiting a session with an open, failed transaction (Tom Lane)
  This bug has no impact on normal builds with asserts not enabled.
• Fix dangling pointer after \texttt{CREATE TABLE AS/SELECT INTO} in a SQL-language function (Tom Lane)
  In most cases this only led to an assertion failure in assert-enabled builds, but worse consequences seem possible.
• Avoid double close of file handle in syslogger on Windows (MauMau)
  Ordinarily this error was invisible, but it would cause an exception when running on a debug version of Windows.
• Fix I/O-conversion-related memory leaks in \texttt{plpgsql} (Andres Freund, Jan Urbanski, Tom Lane)
  Certain operations would leak memory until the end of the current function.
• Work around bug in perl's \texttt{SvPVutf8()} function (Andrew Dunstan)
  This function crashes when handed a typeglob or certain read-only objects such as \texttt{$^V}. Make \texttt{plperl} avoid passing those to it.
• In \texttt{pg_dump}, don't dump contents of an extension's configuration tables if the extension itself is not being dumped (Tom Lane)
• Improve \texttt{pg_dump}'s handling of inherited table columns (Tom Lane)
  \texttt{pg_dump} mishandled situations where a child column has a different default expression than its parent column. If the default is textually identical to the parent's default, but not actually the same (for instance, because of schema search path differences) it would not be recognized as different, so that after dump and restore the child would be allowed to inherit the parent's default. Child columns that are \texttt{NOT NULL} where their parent is not could also be restored subtly incorrectly.
• Fix \texttt{pg_restore}'s direct-to-database mode for \texttt{INSERT}-style table data (Tom Lane)
  Direct-to-database restores from archive files made with \texttt{--inserts} or \texttt{--column-inserts} options fail when using \texttt{pg_restore} from a release dated September or December 2011, as a result of an oversight in a fix for another problem. The archive file itself is not at fault, and text-mode output is okay.
• Teach \texttt{pg_upgrade} to handle renaming of \texttt{plpython}'s shared library (Bruce Momjian)
  Upgrading a pre-9.1 database that included \texttt{plpython} would fail because of this oversight.
• Allow \texttt{pg_upgrade} to process tables containing \texttt{regclass} columns (Bruce Momjian)
  Since \texttt{pg_upgrade} now takes care to preserve \texttt{pg_class} OIDs, there was no longer any reason for this restriction.
• Make libpq ignore ENOTDIR errors when looking for an SSL client certificate file (Magnus Hagander)

  This allows SSL connections to be established, though without a certificate, even when the user's home directory is set to something like /dev/null.

• Fix some more field alignment issues in ecpg's SQLDA area (Zoltan Boszormenyi)

• Allow AT option in ecpg DEALLOCATE statements (Michael Meskes)

  The infrastructure to support this has been there for awhile, but through an oversight there was still an error check rejecting the case.

• Do not use the variable name when defining a varchar structure in ecpg (Michael Meskes)

• Fix contrib/auto_explain's JSON output mode to produce valid JSON (Andrew Dunstan)

  The output used brackets at the top level, when it should have used braces.

• Fix error in contrib/intarray's int[] & int[] operator (Guillaume Lelarge)

  If the smallest integer the two input arrays have in common is 1, and there are smaller values in either array, then 1 would be incorrectly omitted from the result.

• Fix error detection in contrib/pgcrypto's encrypt_iv() and decrypt_iv() (Marko Kreen)

  These functions failed to report certain types of invalid-input errors, and would instead return random garbage values for incorrect input.

• Fix one-byte buffer overrun in contrib/test_parser (Paul Guyot)

  The code would try to read one more byte than it should, which would crash in corner cases. Since contrib/test_parser is only example code, this is not a security issue in itself, but bad example code is still bad.

• Use __sync_lock_test_and_set() for spinlocks on ARM, if available (Martin Pitt)

  This function replaces our previous use of the SWPB instruction, which is deprecated and not available on ARMv6 and later. Reports suggest that the old code doesn't fail in an obvious way on recent ARM boards, but simply doesn't interlock concurrent accesses, leading to bizarre failures in multiprocess operation.

• Use -fexcess-precision=standard option when building with gcc versions that accept it (Andrew Dunstan)

  This prevents assorted scenarios wherein recent versions of gcc will produce creative results.

• Allow use of threaded Python on FreeBSD (Chris Rees)

  Our configure script previously believed that this combination wouldn't work; but FreeBSD fixed the problem, so remove that error check.

• Allow MinGW builds to use standardly-named OpenSSL libraries (Tomasz Ostrowski)

### E.132. Release 9.1.2

**Release date:** 2011-12-05

This release contains a variety of fixes from 9.1.1. For information about new features in the 9.1 major release, see Section E.134.
E.132.1. Migration to Version 9.1.2

A dump/restore is not required for those running 9.1.X.

However, a longstanding error was discovered in the definition of the information_schema.referential_constraints view. If you rely on correct results from that view, you should replace its definition as explained in the first changelog item below.

Also, if you use the citext data type, and you upgraded from a previous major release by running pg_upgrade, you should run CREATE EXTENSION citext FROM unpackaged to avoid collation-related failures in citext operations. The same is necessary if you restore a dump from a pre-9.1 database that contains an instance of the citext data type. If you've already run the CREATE EXTENSION command before upgrading to 9.1.2, you will instead need to do manual catalog updates as explained in the second changelog item.

E.132.2. Changes

• Fix bugs in information_schema.referential_constraints view (Tom Lane)

  This view was being insufficiently careful about matching the foreign-key constraint to the depend-ed-on primary or unique key constraint. That could result in failure to show a foreign key constraint at all, or showing it multiple times, or claiming that it depends on a different constraint than the one it really does.

  Since the view definition is installed by initdb, merely upgrading will not fix the prob-lem. If you need to fix this in an existing installation, you can (as a superuser) drop the information_schema schema then re-create it by sourcing $SHAREDIR/informa-tion_schema.sql. (Run pg_config --sharedir if you're uncertain where $SHAREDIR is.) This must be repeated in each database to be fixed.

• Make contrib/citext's upgrade script fix collations of citext columns and indexes (Tom Lane)

  Existing citext columns and indexes aren't correctly marked as being of a collatable data type during pg_upgrade from a pre-9.1 server, or when a pre-9.1 dump containing the citext type is loaded into a 9.1 server. That leads to operations on these columns failing with errors such as “could not determine which collation to use for string comparison”. This change allows them to be fixed by the same script that upgrades the citext module into a proper 9.1 extension during CREATE EXTENSION citext FROM unpackaged.

  If you have a previously-upgraded database that is suffering from this problem, and you already ran the CREATE EXTENSION command, you can manually run (as superuser) the UPDATE commands found at the end of $SHAREDIR/extension/citext--unpackaged--1.0.sql. (Run pg_config --sharedir if you're uncertain where $SHAREDIR is.) There is no harm in doing this again if unsure.

• Fix possible crash during UPDATE or DELETE that joins to the output of a scalar-returning function (Tom Lane)

  A crash could only occur if the target row had been concurrently updated, so this problem surfaced only intermittently.

• Fix incorrect replay of WAL records for GIN index updates (Tom Lane)

  This could result in transiently failing to find index entries after a crash, or on a hot-standby server. The problem would be repaired by the next VACUUM of the index, however.

• Fix TOAST-related data corruption during CREATE TABLE dest AS SELECT * FROM src or INSERT INTO dest SELECT * FROM src (Tom Lane)
If a table has been modified by `ALTER TABLE ADD COLUMN`, attempts to copy its data verbatim
to another table could produce corrupt results in certain corner cases. The problem can only manifest
in this precise form in 8.4 and later, but we patched earlier versions as well in case there are other
code paths that could trigger the same bug.

• Fix possible failures during hot standby startup (Simon Riggs)

• Start hot standby faster when initial snapshot is incomplete (Simon Riggs)

• Fix race condition during toast table access from stale syscache entries (Tom Lane)

The typical symptom was transient errors like “missing chunk number 0 for toast value NNNNN in
pg_toast_2619”, where the cited toast table would always belong to a system catalog.

• Track dependencies of functions on items used in parameter default expressions (Tom Lane)

Previously, a referenced object could be dropped without having dropped or modified the function,
leading to misbehavior when the function was used. Note that merely installing this update will not
fix the missing dependency entries; to do that, you’d need to `CREATE OR REPLACE` each such
function afterwards. If you have functions whose defaults depend on non-built-in objects, doing so
is recommended.

• Fix incorrect management of placeholder variables in nestloop joins (Tom Lane)

This bug is known to lead to “variable not found in subplan target list” planner errors, and could
possibly result in wrong query output when outer joins are involved.

• Fix window functions that sort by expressions involving aggregates (Tom Lane)

Previously these could fail with “could not find pathkey item to sort” planner errors.

• Fix “MergeAppend child’s targetlist doesn’t match MergeAppend” planner errors (Tom Lane)

• Fix index matching for operators with both collatable and noncollatable inputs (Tom Lane)

In 9.1.0, an indexable operator that has a non-collatable left-hand input type and a collatable right-
hand input type would not be recognized as matching the left-hand column’s index. An example is
the `hstore ? text` operator.

• Allow inlining of set-returning SQL functions with multiple OUT parameters (Tom Lane)

• Don’t trust deferred-unique indexes for join removal (Tom Lane and Marti Raudsepp)

A deferred uniqueness constraint might not hold intra-transaction, so assuming that it does could
give incorrect query results.

• Make `DatumGetInetP()` unpack inet datums that have a 1-byte header, and add a new macro,
`DatumGetInetPP()`, that does not (Heikki Linnakangas)

This change affects no core code, but might prevent crashes in add-on code that expects`DatumGetInetP()`
to produce an unpacked datum as per usual convention.

• Improve locale support in `money` type’s input and output (Tom Lane)

Aside from not supporting all standard `lc_monetary` formatting options, the input and output
functions were inconsistent, meaning there were locales in which dumped `money` values could not
be re-read.

• `transform_null_equals` affects `CASE` foo WHEN NULL ... constructs (Heikki
Linnakangas)
transform_null_equals is only supposed to affect \texttt{foo} = \texttt{NULL} expressions written directly by the user, not equality checks generated internally by this form of \texttt{CASE}.

- Change foreign-key trigger creation order to better support self-referential foreign keys (Tom Lane)

  For a cascading foreign key that references its own table, a row update will fire both the \texttt{ON UPDATE} trigger and the \texttt{CHECK} trigger as one event. The \texttt{ON UPDATE} trigger must execute first, else the \texttt{CHECK} will check a non-final state of the row and possibly throw an inappropriate error. However, the firing order of these triggers is determined by their names, which generally sort in creation order since the triggers have auto-generated names following the convention “RI\_ConstraintTrigger\_NNNN”. A proper fix would require modifying that convention, which we will do in 9.2, but it seems risky to change it in existing releases. So this patch just changes the creation order of the triggers. Users encountering this type of error should drop and re-create the foreign key constraint to get its triggers into the right order.

- Fix \texttt{IF EXISTS} to work correctly in \texttt{DROP OPERATOR FAMILY} (Robert Haas)

- Disallow dropping of an extension from within its own script (Tom Lane)

  This prevents odd behavior in case of incorrect management of extension dependencies.

- Don’t mark auto-generated types as extension members (Robert Haas)

  Relation rowtypes and automatically-generated array types do not need to have their own extension membership entries in \texttt{pg\_depend}, and creating such entries complicates matters for extension upgrades.

- Cope with invalid pre-existing \texttt{search\_path} settings during \texttt{CREATE EXTENSION} (Tom Lane)

- Avoid floating-point underflow while tracking buffer allocation rate (Greg Matthews)

  While harmless in itself, on certain platforms this would result in annoying kernel log messages.

- Prevent autovacuum transactions from running in serializable mode (Tom Lane)

  Autovacuum formerly used the cluster-wide default transaction isolation level, but there is no need for it to use anything higher than \texttt{READ COMMITTED}, and using \texttt{SERIALIZABLE} could result in unnecessary delays for other processes.

- Ensure walsender processes respond promptly to \texttt{SIGTERM} (Magnus Hagander)

- Exclude \texttt{postmaster.opts} from base backups (Magnus Hagander)

- Preserve configuration file name and line number values when starting child processes under Windows (Tom Lane)

  Formerly, these would not be displayed correctly in the \texttt{pg\_settings} view.

- Fix incorrect field alignment in ecpg’s SQLDA area (Zoltan Boszormenyi)

- Preserve blank lines within commands in psql’s command history (Robert Haas)

  The former behavior could cause problems if an empty line was removed from within a string literal, for example.

- Avoid platform-specific infinite loop in \texttt{pg\_dump} (Steve Singer)

- Fix compression of plain-text output format in \texttt{pg\_dump} (Adrian Klaver and Tom Lane)

  \texttt{pg\_dump} has historically understood \texttt{-Z} with no \texttt{-F} switch to mean that it should emit a gzip-compressed version of its plain text output. Restore that behavior.
• Fix pg_dump to dump user-defined casts between auto-generated types, such as table rowtypes (Tom Lane)

• Fix missed quoting of foreign server names in pg_dump (Tom Lane)

• Assorted fixes for pg_upgrade (Bruce Momjian)

  Handle exclusion constraints correctly, avoid failures on Windows, don't complain about mismatched toast table names in 8.4 databases.

• In PL/pgSQL, allow foreign tables to define row types (Alexander Soudakov)

• Fix up conversions of PL/Perl functions’ results (Alex Hunsaker and Tom Lane)

  Restore the pre-9.1 behavior that PL/Perl functions returning void ignore the result value of their last Perl statement; 9.1.0 would throw an error if that statement returned a reference. Also, make sure it works to return a string value for a composite type, so long as the string meets the type’s input format. In addition, throw errors for attempts to return Perl arrays or hashes when the function’s declared result type is not an array or composite type, respectively. (Pre-9.1 versions rather uselessly returned strings like ARRAY(0x221a9a0) or HASH(0x221aa90) in such cases.)

• Ensure PL/Perl strings are always correctly UTF8-encoded (Amit Khandekar and Alex Hunsaker)

• Use the preferred version of xsubpp to build PL/Perl, not necessarily the operating system’s main copy (David Wheeler and Alex Hunsaker)

• Correctly propagate SQLSTATE in PL/Python exceptions (Mika Eloranta and Jan Urbanski)

• Do not install PL/Python extension files for Python major versions other than the one built against (Peter Eisentraut)

• Change all the contrib extension script files to report a useful error message if they are fed to psql (Andrew Dunstan and Tom Lane)

  This should help teach people about the new method of using CREATE EXTENSION to load these files. In most cases, sourcing the scripts directly would fail anyway, but with harder-to-interpret messages.

• Fix incorrect coding in contrib/dict_int and contrib/dict_xsyn (Tom Lane)

  Some functions incorrectly assumed that memory returned by palloc() is guaranteed zeroed.

• Remove contrib/sepgsql tests from the regular regression test mechanism (Tom Lane)

  Since these tests require root privileges for setup, they’re impractical to run automatically. Switch over to a manual approach instead, and provide a testing script to help with that.

• Fix assorted errors in contrib/unaccent’s configuration file parsing (Tom Lane)

• Honor query cancel interrupts promptly in pgstatindex() (Robert Haas)

• Fix incorrect quoting of log file name in macOS start script (Sidar Lopez)

• Revert unintentional enabling of WAL_DEBUG (Robert Haas)

  Fortunately, as debugging tools go, this one is pretty cheap; but it’s not intended to be enabled by default, so revert.

• Ensure VPATH builds properly install all server header files (Peter Eisentraut)

• Shorten file names reported in verbose error messages (Peter Eisentraut)
Regular builds have always reported just the name of the C file containing the error message call, but VPATH builds formerly reported an absolute path name.

- Fix interpretation of Windows timezone names for Central America (Tom Lane)

  Map “Central America Standard Time” to \texttt{CST6}, not \texttt{CST6CDT}, because DST is generally not observed anywhere in Central America.

- Update time zone data files to tzdata release 2011n for DST law changes in Brazil, Cuba, Fiji, Palestine, Russia, and Samoa; also historical corrections for Alaska and British East Africa.

\textbf{E.133. Release 9.1.1}

\textbf{Release date:} 2011-09-26

This release contains a small number of fixes from 9.1.0. For information about new features in the 9.1 major release, see Section E.134.

\textbf{E.133.1. Migration to Version 9.1.1}

A dump/restore is not required for those running 9.1.X.

\textbf{E.133.2. Changes}

- Make \texttt{pg\_options\_to\_table} return NULL for an option with no value (Tom Lane)

  Previously such cases would result in a server crash.

- Fix memory leak at end of a GiST index scan (Tom Lane)

  Commands that perform many separate GiST index scans, such as verification of a new GiST-based exclusion constraint on a table already containing many rows, could transiently require large amounts of memory due to this leak.

- Fix explicit reference to \texttt{pg\_temp} schema in \texttt{CREATE TEMPORARY TABLE} (Robert Haas)

  This used to be allowed, but failed in 9.1.0.

\textbf{E.134. Release 9.1}

\textbf{Release date:} 2011-09-12

\textbf{E.134.1. Overview}

This release shows PostgreSQL moving beyond the traditional relational-database feature set with new, ground-breaking functionality that is unique to PostgreSQL. The streaming replication feature introduced in release 9.0 is significantly enhanced by adding a synchronous-replication option, streaming backups, and monitoring improvements. Major enhancements include:

- Allow synchronous replication
- Add support for foreign tables
- Add per-column collation support
- Add extensions which simplify packaging of additions to PostgreSQL
- Add a true serializable isolation level
• Support unlogged tables using the UNLOGGED option in CREATE TABLE
• Allow data-modification commands (INSERT/UPDATE/DELETE) in WITH clauses
• Add nearest-neighbor (order-by-operator) searching to GiST indexes
• Add a SECURITY LABEL command and support for SELinux permissions control
• Update the PL/Python server-side language

The above items are explained in more detail in the sections below.

E.134.2. Migration to Version 9.1

A dump/restore using pg_dump, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.1 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:

E.134.2.1. Strings

• Change the default value of standard_conforming_strings to on (Robert Haas)

By default, backslashes are now ordinary characters in string literals, not escape characters. This change removes a long-standing incompatibility with the SQL standard. escape_string_warning has produced warnings about this usage for years. E'' strings are the proper way to embed backslash escapes in strings and are unaffected by this change.

Warning

This change can break applications that are not expecting it and do their own string escaping according to the old rules. The consequences could be as severe as introducing SQL-injection security holes. Be sure to test applications that are exposed to untrusted input, to ensure that they correctly handle single quotes and backslashes in text strings.

E.134.2.2. Casting

• Disallow function-style and attribute-style data type casts for composite types (Tom Lane)

For example, disallow composite_value.text and text(composite_value). Unintentional uses of this syntax have frequently resulted in bug reports; although it was not a bug, it seems better to go back to rejecting such expressions. The CAST and :: syntaxes are still available for use when a cast of an entire composite value is actually intended.

• Tighten casting checks for domains based on arrays (Tom Lane)

When a domain is based on an array type, it is allowed to “look through” the domain type to access the array elements, including subscribing the domain value to fetch or assign an element. Assignment to an element of such a domain value, for instance via UPDATE ... SET domaincol[5] = .... will now result in rechecking the domain type's constraints, whereas before the checks were skipped.

E.134.2.3. Arrays

• Change string_to_array() to return an empty array for a zero-length string (Pavel Stehule)
Previously this returned a null value.

- Change `string_to_array()` so a NULL separator splits the string into characters (Pavel Stehule)

Previously this returned a null value.

### E.134.2.4. Object Modification

- Fix improper checks for before/after triggers (Tom Lane)

Triggers can now be fired in three cases: BEFORE, AFTER, or INSTEAD OF some action. Trigger function authors should verify that their logic behaves sanely in all three cases.

- Require superuser or `CREATEROLE` permissions in order to set comments on roles (Tom Lane)

### E.134.2.5. Server Settings

- Change `pg_last_xlog_receive_location()` so it never moves backwards (Fujii Masao)

Previously, the value of `pg_last_xlog_receive_location()` could move backward when streaming replication is restarted.

- Have logging of replication connections honor `log_connections` (Magnus Hagander)

Previously, replication connections were always logged.

### E.134.2.6. PL/pgSQL Server-Side Language

- Change PL/pgSQL’s `RAISE` command without parameters to be catchable by the attached exception block (Piyush Newe)

Previously `RAISE` in a code block was always scoped to an attached exception block, so it was uncatchable at the same scope.

- Adjust PL/pgSQL’s error line numbering code to be consistent with other PLs (Pavel Stehule)

Previously, PL/pgSQL would ignore (not count) an empty line at the start of the function body. Since this was inconsistent with all other languages, the special case was removed.

- Make PL/pgSQL complain about conflicting IN and OUT parameter names (Tom Lane)

Formerly, the collision was not detected, and the name would just silently refer to only the OUT parameter.

- Type modifiers of PL/pgSQL variables are now visible to the SQL parser (Tom Lane)

A type modifier (such as a varchar length limit) attached to a PL/pgSQL variable was formerly enforced during assignments, but was ignored for all other purposes. Such variables will now behave more like table columns declared with the same modifier. This is not expected to make any visible difference in most cases, but it could result in subtle changes for some SQL commands issued by PL/pgSQL functions.

### E.134.2.7. Contrib

- All contrib modules are now installed with `CREATE EXTENSION` rather than by manually invok-

ing their SQL scripts (Dimitri Fontaine, Tom Lane)

To update an existing database containing the 9.0 version of a contrib module, use `CREATE EX-

TENSION ... FROM unpackaged` to wrap the existing contrib module's objects into an exten-
When updating from a pre-9.0 version, drop the contrib module’s objects using its old uninstall script, then use \texttt{CREATE EXTENSION}.

### E.134.2.8. Other Incompatibilities

- Make \texttt{pg\_stat\_reset()} reset all database-level statistics (Tomas Vondra)
  
  Some \texttt{pg\_stat\_database} counters were not being reset.

- Fix some \texttt{information\_schema\_triggers} column names to match the new SQL-standard names (Dean Rasheed)

- Treat ECPG cursor names as case-insensitive (Zoltan Boszormenyi)

### E.134.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.1 and the previous major release.

#### E.134.3.1. Server

##### E.134.3.1.1. Performance

- Support unlogged tables using the \texttt{UNLOGGED} option in \texttt{CREATE TABLE} (Robert Haas)
  
  Such tables provide better update performance than regular tables, but are not crash-safe: their contents are automatically cleared in case of a server crash. Their contents do not propagate to replication slaves, either.

- Allow \texttt{FULL OUTER JOIN} to be implemented as a hash join, and allow either side of a \texttt{LEFT OUTER JOIN} or \texttt{RIGHT OUTER JOIN} to be hashed (Tom Lane)
  
  Previously \texttt{FULL OUTER JOIN} could only be implemented as a merge join, and \texttt{LEFT OUTER JOIN} and \texttt{RIGHT OUTER JOIN} could hash only the nullable side of the join. These changes provide additional query optimization possibilities.

- Merge duplicate fsync requests (Robert Haas, Greg Smith)
  
  This greatly improves performance under heavy write loads.

- Improve performance of \texttt{commit\_siblings} (Greg Smith)
  
  This allows the use of \texttt{commit\_siblings} with less overhead.

- Reduce the memory requirement for large ispell dictionaries (Pavel Stehule, Tom Lane)

- Avoid leaving data files open after “blind writes” (Alvaro Herrera)
  
  This fixes scenarios in which backends might hold files open long after they were deleted, preventing the kernel from reclaiming disk space.

##### E.134.3.1.2. Optimizer

- Allow inheritance table scans to return meaningfully-sorted results (Greg Stark, Hans-Jurgen Schonig, Robert Haas, Tom Lane)
  
  This allows better optimization of queries that use \texttt{ORDER BY}, \texttt{LIMIT}, or \texttt{MIN/\texttt{MAX}} with inherited tables.

- Improve GIN index scan cost estimation (Teodor Sigaev)
• Improve cost estimation for aggregates and window functions (Tom Lane)

E.134.3.1.3. Authentication

• Support host names and host suffixes (e.g., example.com) in pg_hba.conf (Peter Eisentraut)
  Previously only host IP addresses and CIDR values were supported.

• Support the key word all in the host column of pg_hba.conf (Peter Eisentraut)
  Previously people used 0.0.0.0/0 or ::/0 for this.

• Reject local lines in pg_hba.conf on platforms that don’t support Unix-socket connections (Magnus Hagander)
  Formerly, such lines were silently ignored, which could be surprising. This makes the behavior more like other unsupported cases.

• Allow GSSAPI to be used to authenticate to servers via SSPI (Christian Ullrich)
  Specifically this allows Unix-based GSSAPI clients to do SSPI authentication with Windows servers.

• ident authentication over local sockets is now known as peer (Magnus Hagander)
  The old term is still accepted for backward compatibility, but since the two methods are fundamentally different, it seemed better to adopt different names for them.

• Rewrite peer authentication to avoid use of credential control messages (Tom Lane)
  This change makes the peer authentication code simpler and better-performing. However, it requires the platform to provide the getpeerid function or an equivalent socket operation. So far as is known, the only platform for which peer authentication worked before and now will not is pre-5.0 NetBSD.

E.134.3.1.4. Monitoring

• Add details to the logging of restartpoints and checkpoints, which is controlled by log_checkpoints (Fujii Masao, Greg Smith)
  New details include WAL file and sync activity.

• Add log_file_mode which controls the permissions on log files created by the logging collector (Martin Pihlak)

• Reduce the default maximum line length for syslog logging to 900 bytes plus prefixes (Noah Misch)
  This avoids truncation of long log lines on syslog implementations that have a 1KB length limit, rather than the more common 2KB.

E.134.3.1.5. Statistical Views

• Add client_hostname column to pg_stat_activity (Peter Eisentraut)
  Previously only the client address was reported.

• Add pg_stat_xact_* statistics functions and views (Joel Jacobson)
  These are like the database-wide statistics counter views, but reflect counts for only the current transaction.

• Add time of last reset in database-level and background writer statistics views (Tomas Vondra)
• Add columns showing the number of vacuum and analyze operations in pg_stat_*_tables views (Magnus Hagander)

• Add buffers_backend_fsync column to pg_stat_bgwriter (Greg Smith)

This new column counts the number of times a backend fsyncs a buffer.

E.134.3.1.6. Server Settings

• Provide auto-tuning of wal_buffers (Greg Smith)

By default, the value of wal_buffers is now chosen automatically based on the value of shared_buffers.

• Increase the maximum values for deadlock_timeout, log_min_duration_statement, and log_autovacuum_min_duration (Peter Eisentraut)

The maximum value for each of these parameters was previously only about 35 minutes. Much larger values are now allowed.

E.134.3.2. Replication and Recovery

E.134.3.2.1. Streaming Replication and Continuous Archiving

• Allow synchronous replication (Simon Riggs, Fujii Masao)

This allows the primary server to wait for a standby to write a transaction's information to disk before acknowledging the commit. One standby at a time can take the role of the synchronous standby, as controlled by the synchronous_standby_names setting. Synchronous replication can be enabled or disabled on a per-transaction basis using the synchronous_commit setting.

• Add protocol support for sending file system backups to standby servers using the streaming replication network connection (Magnus Hagander, Heikki Linnakangas)

This avoids the requirement of manually transferring a file system backup when setting up a standby server.

• Add replication_timeout setting (Fujii Masao, Heikki Linnakangas)

Replication connections that are idle for more than the replication_timeout interval will be terminated automatically. Formerly, a failed connection was typically not detected until the TCP timeout elapsed, which is inconveniently long in many situations.

• Add command-line tool pg_basebackup for creating a new standby server or database backup (Magnus Hagander)

• Add a replication permission for roles (Magnus Hagander)

This is a read-only permission used for streaming replication. It allows a non-superuser role to be used for replication connections. Previously only superusers could initiate replication connections; superusers still have this permission by default.

E.134.3.2.2. Replication Monitoring

• Add system view pg_stat_replication which displays activity of WAL sender processes (Itagaki Takahiro, Simon Riggs)

This reports the status of all connected standby servers.

• Add monitoring function pg_last_xact_replay_timestamp() (Fujii Masao)
This returns the time at which the primary generated the most recent commit or abort record applied on the standby.

### E.134.3.2.3. Hot Standby

- Add configuration parameter `hot_standby_feedback` to enable standbys to postpone cleanup of old row versions on the primary (Simon Riggs)

  This helps avoid canceling long-running queries on the standby.

- Add the `pg_stat_database_conflicts` system view to show queries that have been canceled and the reason (Magnus Hagander)

  Cancellations can occur because of dropped tablespaces, lock timeouts, old snapshots, pinned buffers, and deadlocks.

- Add a `conflicts` count to `pg_stat_database` (Magnus Hagander)

  This is the number of conflicts that occurred in the database.

- Increase the maximum values for `max_standby_archive_delay` and `max_standby_streaming_delay`

  The maximum value for each of these parameters was previously only about 35 minutes. Much larger values are now allowed.

- Add `ERRCODE_T_R_DATABASE_DROPPED` error code to report recovery conflicts due to dropped databases (Tatsuo Ishii)

  This is useful for connection pooling software.

### E.134.3.2.4. Recovery Control

- Add functions to control streaming replication replay (Simon Riggs)

  The new functions are `pg_xlog_replay_pause()`, `pg_xlog_replay_resume()`, and the status function `pg_is_xlog_replay_paused()`.

- Add `recovery.conf` setting `pause_at_recovery_target` to pause recovery at target (Simon Riggs)

  This allows a recovery server to be queried to check whether the recovery point is the one desired.

- Add the ability to create named restore points using `pg_create_restore_point()` (Jaime Casanova)

  These named restore points can be specified as recovery targets using the new `recovery.conf` setting `recovery_target_name`.

- Allow standby recovery to switch to a new timeline automatically (Heikki Linnakangas)

  Now standby servers scan the archive directory for new timelines periodically.

- Add `restart_after_crash` setting which disables automatic server restart after a backend crash (Robert Haas)

  This allows external cluster management software to control whether the database server restarts or not.

- Allow `recovery.conf` to use the same quoting behavior as `postgresql.conf` (Dimitri Fontaine)
Previously all values had to be quoted.

E.134.3.3. Queries

- Add a true serializable isolation level (Kevin Grittner, Dan Ports)
  Previously, asking for serializable isolation guaranteed only that a single MVCC snapshot would be used for the entire transaction, which allowed certain documented anomalies. The old snapshot isolation behavior is still available by requesting the REPEATABLE READ isolation level.

- Allow data-modification commands (INSERT/UPDATE/DELETE) in WITH clauses (Marko Tiikkaja, Hitoshi Harada)
  These commands can use RETURNING to pass data up to the containing query.

- Allow WITH clauses to be attached to INSERT, UPDATE, DELETE statements (Marko Tiikkaja, Hitoshi Harada)

- Allow non-GROUP BY columns in the query target list when the primary key is specified in the GROUP BY clause (Peter Eisentraut)
  The SQL standard allows this behavior, and because of the primary key, the result is unambiguous.

- Allow use of the key word DISTINCT in UNION/INTERSECT/EXCEPT clauses (Tom Lane)
  DISTINCT is the default behavior so use of this key word is redundant, but the SQL standard allows it.

- Fix ordinary queries with rules to use the same snapshot behavior as EXPLAIN ANALYZE (Marko Tiikkaja)
  Previously EXPLAIN ANALYZE used slightly different snapshot timing for queries involving rules. The EXPLAIN ANALYZE behavior was judged to be more logical.

E.134.3.3.1. Strings

- Add per-column collation support (Peter Eisentraut, Tom Lane)
  Previously collation (the sort ordering of text strings) could only be chosen at database creation. Collation can now be set per column, domain, index, or expression, via the SQL-standard COLLATE clause.

E.134.3.4. Object Manipulation

- Add extensions which simplify packaging of additions to PostgreSQL (Dimitri Fontaine, Tom Lane)
  Extensions are controlled by the new CREATE/ALTER/DROP EXTENSION commands. This replaces ad-hoc methods of grouping objects that are added to a PostgreSQL installation.

- Add support for foreign tables (Shigeru Hanada, Robert Haas, Jan Urbanski, Heikki Linnakangas)
  This allows data stored outside the database to be used like native PostgreSQL-stored data. Foreign tables are currently read-only, however.

- Allow new values to be added to an existing enum type via ALTER TYPE (Andrew Dunstan)

- Add ALTER TYPE ... ADD/DROP/ALTER/RENAME ATTRIBUTE (Peter Eisentraut)
  This allows modification of composite types.
E.134.3.4.1. ALTER Object

- Add RESTRICT/CASCADE to ALTER TYPE operations on typed tables (Peter Eisentraut)
  This controls ADD/DROP/ALTER/RENAME ATTRIBUTE cascading behavior.
- Support ALTER TABLE name {OF | NOT OF} type (Noah Misch)
  This syntax allows a standalone table to be made into a typed table, or a typed table to be made
  standalone.
- Add support for more object types in ALTER ... SET SCHEMA commands (Dimitri Fontaine)
  This command is now supported for conversions, operators, operator classes, operator families, text
  search configurations, text search dictionaries, text search parsers, and text search templates.

E.134.3.4.2. CREATE/ALTER TABLE

- Add ALTER TABLE ... ADD UNIQUE/PRIMARY KEY USING INDEX (Gurjeet Singh)
  This allows a primary key or unique constraint to be defined using an existing unique index, including
  a concurrently created unique index.
- Allow ALTER TABLE to add foreign keys without validation (Simon Riggs)
  The new option is called NOT VALID. The constraint’s state can later be modified to VALIDATED
  and validation checks performed. Together these allow you to add a foreign key with minimal impact
  on read and write operations.
- Allow ALTER TABLE ... SET DATA TYPE to avoid table rewrites in appropriate cases (Noah
  Misch, Robert Haas)
  For example, converting a varchar column to text no longer requires a rewrite of the table.
  However, increasing the length constraint on a varchar column still requires a table rewrite.
- Add CREATE TABLE IF NOT EXISTS syntax (Robert Haas)
  This allows table creation without causing an error if the table already exists.
- Fix possible “tuple concurrently updated” error when two backends attempt to add an inheritance
  child to the same table at the same time (Robert Haas)
  ALTER TABLE now takes a stronger lock on the parent table, so that the sessions cannot try to
  update it simultaneously.

E.134.3.4.3. Object Permissions

- Add a SECURITY LABEL command (KaiGai Kohei)
  This allows security labels to be assigned to objects.

E.134.3.5. Utility Operations

- Add transaction-level advisory locks (Marko Tiikkaja)
  These are similar to the existing session-level advisory locks, but such locks are automatically re-
  leased at transaction end.
- Make TRUNCATE ... RESTART IDENTITY restart sequences transactionally (Steve Singer)
  Previously the counter could have been left out of sync if a backend crashed between the on-commit
  truncation activity and commit completion.
E.134.3.5.1. **COPY**

- Add ENC encoding option to COPY TO/FROM (Hitoshi Harada, Itagaki Takahiro)
  
  This allows the encoding of the COPY file to be specified separately from client encoding.

- Add bidirectional COPY protocol support (Fujii Masao)
  
  This is currently only used by streaming replication.

E.134.3.5.2. **EXPLAIN**

- Make EXPLAIN VERBOSE show the function call expression in a FunctionScan node (Tom Lane)

E.134.3.5.3. **VACUUM**

- Add additional details to the output of VACUUM FULL VERBOSE and CLUSTER VERBOSE (Itagaki Takahiro)
  
  New information includes the live and dead tuple count and whether CLUSTER is using an index to rebuild.

- Prevent autovacuum from waiting if it cannot acquire a table lock (Robert Haas)
  
  It will try to vacuum that table later.

E.134.3.5.4. **CLUSTER**

- Allow CLUSTER to sort the table rather than scanning the index when it seems likely to be cheaper (Leonardo Francalanci)

E.134.3.5.5. **Indexes**

- Add nearest-neighbor (order-by-operator) searching to GiST indexes (Teodor Sigaev, Tom Lane)
  
  This allows GiST indexes to quickly return the $N$ closest values in a query with LIMIT. For example

  ```
  SELECT * FROM places ORDER BY location <-> point '(101,456)' LIMIT 10;
  ```

  finds the ten places closest to a given target point.

- Allow GIN indexes to index null and empty values (Tom Lane)
  
  This allows full GIN index scans, and fixes various corner cases in which GIN scans would fail.

- Allow GIN indexes to better recognize duplicate search entries (Tom Lane)
  
  This reduces the cost of index scans, especially in cases where it avoids unnecessary full index scans.

- Fix GiST indexes to be fully crash-safe (Heikki Linnakangas)
  
  Previously there were rare cases where a REINDEX would be required (you would be informed).

E.134.3.6. **Data Types**

- Allow numeric to use a more compact, two-byte header in common cases (Robert Haas)
Previously all numeric values had four-byte headers; this change saves on disk storage.

- Add support for dividing money by money (Andy Balholm)
- Allow binary I/O on type void (Radoslaw Smogura)
- Improve hypotenuse calculations for geometric operators (Paul Matthews)
  This avoids unnecessary overflows, and may also be more accurate.
- Support hashing array values (Tom Lane)
  This provides additional query optimization possibilities.
- Don’t treat a composite type as sortable unless all its column types are sortable (Tom Lane)
  This avoids possible “could not identify a comparison function” failures at runtime, if it is possible to implement the query without sorting. Also, ANALYZE won’t try to use inappropriate statistics-gathering methods for columns of such composite types.

**E.134.3.6.1. Casting**

- Add support for casting between money and numeric (Andy Balholm)
- Add support for casting from int4 and int8 to money (Joey Adams)
- Allow casting a table's row type to the table's supertype if it's a typed table (Peter Eisentraut)
  This is analogous to the existing facility that allows casting a row type to a supertable's row type.

**E.134.3.6.2. XML**

- Add XML function XMLEXISTS and xpath_exists() functions (Mike Fowler)
  These are used for XPath matching.
- Add XML functions xml_is_well_formed(), xml_is_well_formed_document(), xml_is_well_formed_content() (Mike Fowler)
  These check whether the input is properly-formed XML. They provide functionality that was previously available only in the deprecated contrib/xml2 module.

**E.134.3.7. Functions**

- Add SQL function format(text, ...), which behaves analogously to C's printf() (Pavel Stehule, Robert Haas)
  It currently supports formats for strings, SQL literals, and SQL identifiers.
- Add string functions concat(), concat_ws(), left(), right(), and reverse() (Pavel Stehule)
  These improve compatibility with other database products.
- Add function pg_read_binary_file() to read binary files (Dimitri Fontaine, Itagaki Takahiro)
- Add a single-parameter version of function pg_read_file() to read an entire file (Dimitri Fontaine, Itagaki Takahiro)
- Add three-parameter forms of array_to_string() and string_to_array() for null value processing control (Pavel Stehule)
E.134.3.7.1. Object Information Functions

- Add the `pg_describe_object()` function (Alvaro Herrera)
  This function is used to obtain a human-readable string describing an object, based on the `pg_class` OID, object OID, and sub-object ID. It can be used to help interpret the contents of `pg_depend`.
- Update comments for built-in operators and their underlying functions (Tom Lane)
  Functions that are meant to be used via an associated operator are now commented as such.
- Add variable `quote_all_identifiers` to force the quoting of all identifiers in `EXPLAIN` and in system catalog functions like `pg_get_viewdef()` (Robert Haas)
  This makes exporting schemas to tools and other databases with different quoting rules easier.
- Add columns to the `information_schema.sequences` system view (Peter Eisentraut)
  Previously, though the view existed, the columns about the sequence parameters were unimplemented.
- Allow `public` as a pseudo-role name in `has_table_privilege()` and related functions (Alvaro Herrera)
  This allows checking for public permissions.

E.134.3.7.2. Function and Trigger Creation

- Support `INSTEAD OF` triggers on views (Dean Rasheed)
  This feature can be used to implement fully updatable views.

E.134.3.8. Server-Side Languages

E.134.3.8.1. PL/pgSQL Server-Side Language

- Add `FOREACH IN ARRAY` to PL/pgSQL (Pavel Stehule)
  This is more efficient and readable than previous methods of iterating through the elements of an array value.
- Allow `RAISE` without parameters to be caught in the same places that could catch a `RAISE ERROR` from the same location (Piyush Newe)
  The previous coding threw the error from the block containing the active exception handler. The new behavior is more consistent with other DBMS products.

E.134.3.8.2. PL/Perl Server-Side Language

- Allow generic record arguments to PL/Perl functions (Andrew Dunstan)
  PL/Perl functions can now be declared to accept type `record`. The behavior is the same as for any named composite type.
- Convert PL/Perl array arguments to Perl arrays (Alexey Klyukin, Alex Hunsaker)
  String representations are still available.
- Convert PL/Perl composite-type arguments to Perl hashes (Alexey Klyukin, Alex Hunsaker)
  String representations are still available.
E.134.3.8.3. PL/Python Server-Side Language

- Add table function support for PL/Python (Jan Urbanski)
  PL/Python can now return multiple OUT parameters and record sets.
- Add a validator to PL/Python (Jan Urbanski)
  This allows PL/Python functions to be syntax-checked at function creation time.
- Allow exceptions for SQL queries in PL/Python (Jan Urbanski)
  This allows access to SQL-generated exception error codes from PL/Python exception blocks.
- Add explicit subtransactions to PL/Python (Jan Urbanski)
- Add PL/Python functions for quoting strings (Jan Urbanski)
  These functions are `plpy.quote_ident`, `plpy.quote_literal`, and `plpy.quote_nullable`.
- Add traceback information to PL/Python errors (Jan Urbanski)
- Report PL/Python errors from iterators with `PLY_elog` (Jan Urbanski)
- Fix exception handling with Python 3 (Jan Urbanski)
  Exception classes were previously not available in `plpy` under Python 3.

E.134.3.9. Client Applications

- Mark createlang and droplang as deprecated now that they just invoke extension commands (Tom Lane)

E.134.3.9.1. psql

- Add psql command `\conninfo` to show current connection information (David Christensen)
- Add psql command `\sf` to show a function's definition (Pavel Stehule)
- Add psql command `\dll` to list languages (Fernando Ike)
- Add the `S` ("system") option to psql's `\d` (list schemas) command (Tom Lane)
  `\d` without `S` now suppresses system schemas.
- Allow psql's `\e` and `\ef` commands to accept a line number to be used to position the cursor in
  the editor (Pavel Stehule)
  This is passed to the editor according to the `PSQL_EDITOR_LINENUMBER_ARG` environment variable.
- Have psql set the client encoding from the operating system locale by default (Heikki Linnakangas)
  This only happens if the `PGCLIENTENCODING` environment variable is not set.
- Make `\d` distinguish between unique indexes and unique constraints (Josh Kupershmidt)
- Make `\dt+` report `pg_table_size` instead of `pg_relation_size` when talking to 9.0 or later servers (Bernd Helmle)
  This is a more useful measure of table size, but note that it is not identical to what was previously
  reported in the same display.
• Additional tab completion support (Itagaki Takahiro, Pavel Stehule, Andrey Popp, Christoph Berg, David Fetter, Josh Kupershmidt)

E.134.3.9.2. pg_dump

• Add pg_dump and pg_dumpall option --quote-all-identifiers to force quoting of all identifiers (Robert Haas)

• Add directory format to pg_dump (Joachim Wieland, Heikki Linnakangas)
  This is internally similar to the tar pg_dump format.

E.134.3.9.3. pg_ctl

• Fix pg_ctl so it no longer incorrectly reports that the server is not running (Bruce Momjian)
  Previously this could happen if the server was running but pg_ctl could not authenticate.

• Improve pg_ctl start's “wait” (-w) option (Bruce Momjian, Tom Lane)
  The wait mode is now significantly more robust. It will not get confused by non-default postmaster port numbers, non-default Unix-domain socket locations, permission problems, or stale postmaster lock files.

• Add promote option to pg_ctl to switch a standby server to primary (Fujii Masao)

E.134.3.10. Development Tools

E.134.3.10.1. libpq

• Add a libpq connection option client_encoding which behaves like the PGCLIENTENCODING environment variable (Heikki Linnakangas)
  The value auto sets the client encoding based on the operating system locale.

• Add PQlibVersion() function which returns the libpq library version (Magnus Hagander)
  libpq already had PQserverVersion() which returns the server version.

• Allow libpq-using clients to check the user name of the server process when connecting via Unix-domain sockets, with the new requirepeer connection option (Peter Eisentraut)
  PostgreSQL already allowed servers to check the client user name when connecting via Unix-domain sockets.

• Add PQping() and PQpingParams() to libpq (Bruce Momjian, Tom Lane)
  These functions allow detection of the server’s status without trying to open a new session.

E.134.3.10.2. ECPG

• Allow ECPG to accept dynamic cursor names even in WHERE CURRENT OF clauses (Zoltan Boszormenyi)

• Make ecpglib write double values with a precision of 15 digits, not 14 as formerly (Akira Kurosawa)

E.134.3.11. Build Options

• Use +Olibmerrno compile flag with HP-UX C compilers that accept it (Ibrar Ahmed)
  This avoids possible misbehavior of math library calls on recent HP platforms.
E.134.3.11.1. Makefiles

- Improved parallel make support (Peter Eisentraut)
  This allows for faster compiles. Also, `make -k` now works more consistently.
- Require GNU make 3.80 or newer (Peter Eisentraut)
  This is necessary because of the parallel-make improvements.
- Add `make maintainer-check` target (Peter Eisentraut)
  This target performs various source code checks that are not appropriate for either the build or the regression tests. Currently: duplicate_oids, SGML syntax and tabs check, NLS syntax check.
- Support `make check` in `contrib` (Peter Eisentraut)
  Formerly only `make installcheck` worked, but now there is support for testing in a temporary installation. The top-level `make check-world` target now includes testing `contrib` this way.

E.134.3.11.2. Windows

- On Windows, allow pg_ctl to register the service as auto-start or start-on-demand (Quan Zongliang)
- Add support for collecting crash dumps on Windows (Craig Ringer, Magnus Hagander)
  Minidumps can now be generated by non-debug Windows binaries and analyzed by standard debugging tools.
- Enable building with the MinGW64 compiler (Andrew Dunstan)
  This allows building 64-bit Windows binaries even on non-Windows platforms via cross-compiling.

E.134.3.12. Source Code

- Revise the API for GUC variable assign hooks (Tom Lane)
  The previous functions of assign hooks are now split between check hooks and assign hooks, where the former can fail but the latter shouldn't. This change will impact add-on modules that define custom GUC parameters.
- Add latches to the source code to support waiting for events (Heikki Linnakangas)
- Centralize data modification permissions-checking logic (KaiGai Kohei)
- Add missing `get_object_oid()` functions, for consistency (Robert Haas)
- Improve ability to use C++ compilers for compiling add-on modules by removing conflicting key words (Tom Lane)
- Add support for DragonFly BSD (Rumko)
- Expose `quote_literal_cstr()` for backend use (Robert Haas)
- Run regression tests in the default encoding (Peter Eisentraut)
  Regression tests were previously always run with `SQL_ASCII` encoding.
- Add `src/tools/git_changelog` to replace `cvs2cl` and `pgcvslog` (Robert Haas, Tom Lane)
- Add `git-external-diff` script to `src/tools` (Bruce Momjian)
Release Notes

This is used to generate context diffs from git.

- Improve support for building with Clang (Peter Eisentraut)

**E.134.3.12.1. Server Hooks**

- Add source code hooks to check permissions (Robert Haas, Stephen Frost)
- Add post-object-creation function hooks for use by security frameworks (KaiGai Kohei)
- Add a client authentication hook (KaiGai Kohei)

**E.134.3.13. Contrib**

- Modify contrib modules and procedural languages to install via the new extension mechanism (Tom Lane, Dimitri Fontaine)
- Add contrib/file_fdw foreign-data wrapper (Shigeru Hanada)
  
  Foreign tables using this foreign data wrapper can read flat files in a manner very similar to COPY.
- Add nearest-neighbor search support to contrib/pg_trgm and contrib/btree_gist (Teodor Sigaev)
- Add contrib/btree_gist support for searching on not-equals (Jeff Davis)
- Fix contrib/fuzzystromatch's levenshtein() function to handle multibyte characters (Alexander Korotkov)
- Add ssl_cipher() and ssl_version() functions to contrib/sslinfo (Robert Haas)
- Fix contrib/intarray and contrib/hstore to give consistent results with indexed empty arrays (Tom Lane)
  
  Previously an empty-array query that used an index might return different results from one that used a sequential scan.
- Allow contrib/intarray to work properly on multidimensional arrays (Tom Lane)
- In contrib/intarray, avoid errors complaining about the presence of nulls in cases where no nulls are actually present (Tom Lane)
- In contrib/intarray, fix behavior of containment operators with respect to empty arrays (Tom Lane)
  
  Empty arrays are now correctly considered to be contained in any other array.
- Remove contrib/xml2's arbitrary limit on the number of parameter=value pairs that can be handled by xslt_process() (Pavel Stehule)
  
  The previous limit was 10.
- In contrib/pageinspect, fix heap_page_item to return infomasks as 32-bit values (Alvaro Herrera)
  
  This avoids returning negative values, which was confusing. The underlying value is a 16-bit unsigned integer.

**E.134.3.13.1. Security**

- Add contrib/sepgsql to interface permission checks with SELinux (KaiGai Kohei)
This uses the new SECURITY LABEL facility.

- **Add contrib module auth_delay (KaiGai Kohei)**
  This causes the server to pause before returning authentication failure; it is designed to make brute force password attacks more difficult.

- **Add dummy_seclabel contrib module (KaiGai Kohei)**
  This is used for permission regression testing.

### E.134.3.13.2. Performance

- **Add support for LIKE and ILIKE index searches to contrib/pg_trgm (Alexander Korotkov)**

- **Add levenshtein_less_equal() function to contrib/fuzzystrmatch, which is optimized for small distances (Alexander Korotkov)**

- **Improve performance of index lookups on contrib/seg columns (Alexander Korotkov)**

- **Improve performance of pg_upgrade for databases with many relations (Bruce Momjian)**

- **Add flag to contrib/pgbench to report per-statement latencies (Florian Pflug)**

### E.134.3.13.3. Fsync Testing

- **Move src/tools/test_fsync to contrib/pg_test_fsync (Bruce Momjian, Tom Lane)**

- **Add O_DIRECT support to contrib/pg_test_fsync (Bruce Momjian)**
  This matches the use of O_DIRECT by wal_sync_method.

- **Add new tests to contrib/pg_test_fsync (Bruce Momjian)**

### E.134.3.14. Documentation

- **Extensive ECPG documentation improvements (Satoshi Nagayasu)**

- **Extensive proofreading and documentation improvements (Thom Brown, Josh Kupershmidt, Susanne Ebrecht)**

- **Add documentation for exit_on_error (Robert Haas)**
  This parameter causes sessions to exit on any error.

- **Add documentation for pg_options_to_table() (Josh Berkus)**
  This function shows table storage options in a readable form.

- **Document that it is possible to access all composite type fields using (compositeval).* syntax (Peter Eisentraut)**

- **Document that translate() removes characters in from that don't have a corresponding to character (Josh Kupershmidt)**

- **Merge documentation for CREATE CONSTRAINT TRIGGER and CREATE TRIGGER (Alvaro Herrera)**

- **Centralize permission and upgrade documentation (Bruce Momjian)**

- **Add kernel tuning documentation for Solaris 10 (Josh Berkus)**
Previously only Solaris 9 kernel tuning was documented.

- Handle non-ASCII characters consistently in HISTORY file (Peter Eisenraut)

While the HISTORY file is in English, we do have to deal with non-ASCII letters in contributor names. These are now transliterated so that they are reasonably legible without assumptions about character set.

**E.135. Release 9.0.23**

**Release date:** 2015-10-08

This release contains a variety of fixes from 9.0.22. For information about new features in the 9.0 major release, see Section E.158.

This is expected to be the last PostgreSQL release in the 9.0.X series. Users are encouraged to update to a newer release branch soon.

**E.135.1. Migration to Version 9.0.23**

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.18, see Section E.140.

**E.135.2. Changes**

- **Fix contrib/pgcrypto to detect and report too-short crypt() salts (Josh Kupershmidt)**

  Certain invalid salt arguments crashed the server or disclosed a few bytes of server memory. We have not ruled out the viability of attacks that arrange for presence of confidential information in the disclosed bytes, but they seem unlikely. (CVE-2015-5288)

- **Fix subtransaction cleanup after a portal (cursor) belonging to an outer subtransaction fails (Tom Lane, Michael Paquier)**

  A function executed in an outer-subtransaction cursor could cause an assertion failure or crash by referencing a relation created within an inner subtransaction.

- **Fix insertion of relations into the relation cache “init file” (Tom Lane)**

  An oversight in a patch in the most recent minor releases caused pg_trigger_tgrelid_tgname_index to be omitted from the init file. Subsequent sessions detected this, then deemed the init file to be broken and silently ignored it, resulting in a significant degradation in session startup time. In addition to fixing the bug, install some guards so that any similar future mistake will be more obvious.

- **Avoid O(N^2) behavior when inserting many tuples into a SPI query result (Neil Conway)**

- **Improve LISTEN startup time when there are many unread notifications (Matt Newell)**

- **Disable SSL renegotiation by default (Michael Paquier, Andres Freund)**

  While use of SSL renegotiation is a good idea in theory, we have seen too many bugs in practice, both in the underlying OpenSSL library and in our usage of it. Renegotiation will be removed entirely in 9.5 and later. In the older branches, just change the default value of ssl_renegotiation_limit to zero (disabled).

- **Lower the minimum values of the *_freeze_max_age parameters (Andres Freund)**
This is mainly to make tests of related behavior less time-consuming, but it may also be of value for installations with limited disk space.

- Limit the maximum value of `wal_buffers` to 2GB to avoid server crashes (Josh Berkus)
- Fix rare internal overflow in multiplication of numeric values (Dean Rasheed)
- Guard against hard-to-reach stack overflows involving record types, range types, `json`, `jsonb`, `tsquery`, `ltxtquery` and `query_int` (Noah Misch)
- Fix handling of `DOW` and `DOY` in datetime input (Greg Stark)

These tokens aren't meant to be used in datetime values, but previously they resulted in opaque internal error messages rather than “invalid input syntax”.

- Add more query-cancel checks to regular expression matching (Tom Lane)
- Add recursion depth protections to regular expression, `SIMILAR TO`, and `LIKE` matching (Tom Lane)

Suitable search patterns and a low stack depth limit could lead to stack-overrun crashes.

- Fix potential infinite loop in regular expression execution (Tom Lane)

A search pattern that can apparently match a zero-length string, but actually doesn't match because of a back reference, could lead to an infinite loop.

- Fix low-memory failures in regular expression compilation (Andreas Seltenreich)
- Fix low-probability memory leak during regular expression execution (Tom Lane)
- Fix rare low-memory failure in lock cleanup during transaction abort (Tom Lane)
- Fix “unexpected out-of-memory situation during sort” errors when using tuplestores with small `work_mem` settings (Tom Lane)
- Fix very-low-probability stack overrun in `qsort` (Tom Lane)
- Fix “invalid memory alloc request size” failure in hash joins with large `work_mem` settings (Tomas Vondra, Tom Lane)
- Fix assorted planner bugs (Tom Lane)

These mistakes could lead to incorrect query plans that would give wrong answers, or to assertion failures in assert-enabled builds, or to odd planner errors such as “could not devise a query plan for the given query”, “could not find pathkey item to sort”, “plan should not reference subplan's variable”, or “failed to assign all NestLoopParams to plan nodes”. Thanks are due to Andreas Seltenreich and Piotr Stefaniak for fuzz testing that exposed these problems.

- Use fuzzy path cost tiebreaking rule in all supported branches (Tom Lane)

This change is meant to avoid platform-specific behavior when alternative plan choices have effectively-identical estimated costs.

- During postmaster shutdown, ensure that per-socket lock files are removed and listen sockets are closed before we remove the `postmaster.pid` file (Tom Lane)

This avoids race-condition failures if an external script attempts to start a new postmaster as soon as `pg_ctl stop` returns.

- Fix postmaster's handling of a startup-process crash during crash recovery (Tom Lane)
If, during a crash recovery cycle, the startup process crashes without having restored database consistency, we'd try to launch a new startup process, which typically would just crash again, leading to an infinite loop.

- Do not print a WARNING when an autovacuum worker is already gone when we attempt to signal it, and reduce log verbosity for such signals (Tom Lane)
- Prevent autovacuum launcher from sleeping unduly long if the server clock is moved backwards a large amount (Álvaro Herrera)
- Ensure that cleanup of a GIN index's pending-insertions list is interruptable by cancel requests (Jeff Janes)
- Allow all-zeroes pages in GIN indexes to be reused (Heikki Linnakangas)

Such a page might be left behind after a crash.

- Fix off-by-one error that led to otherwise-harmless warnings about "apparent wraparound" in subtrans/multixact truncation (Thomas Munro)
- Fix misreporting of CONTINUE and MOVE statement types in PL/pgSQL's error context messages (Pavel Stehule, Tom Lane)
- Fix some places in PL/Tcl that neglected to check for failure of malloc() calls (Michael Paquier, Álvaro Herrera)
- Improve libpq's handling of out-of-memory conditions (Michael Paquier, Heikki Linnakangas)
- Fix memory leaks and missing out-of-memory checks in ecpg (Michael Paquier)
- Fix psql's code for locale-aware formatting of numeric output (Tom Lane)

The formatting code invoked by \pset numericlocale on did the wrong thing for some uncommon cases such as numbers with an exponent but no decimal point. It could also mangle already-localized output from the money data type.

- Prevent crash in psql's \c command when there is no current connection (Noah Misch)
- Ensure that temporary files created during a pg_dump run with tar-format output are not world-readable (Michael Paquier)
- Fix pg_dump and pg_upgrade to support cases where the postgres or template1 database is in a non-default tablespace (Marti Raudsepp, Bruce Momjian)
- Fix pg_dump to handle object privileges sanely when dumping from a server too old to have a particular privilege type (Tom Lane)

When dumping functions or procedural languages from pre-7.3 servers, pg_dump would produce GRANT/REVOKE commands that revoked the owner's grantable privileges and instead granted all privileges to PUBLIC. Since the privileges involved are just USAGE and EXECUTE, this isn't a security problem, but it's certainly a surprising representation of the older systems' behavior. Fix it to leave the default privilege state alone in these cases.

- Fix pg_dump to dump shell types (Tom Lane)

Shell types (that is, not-yet-fully-defined types) aren't useful for much, but nonetheless pg_dump should dump them.

- Fix spinlock assembly code for PPC hardware to be compatible with AIX's native assembler (Tom Lane)
Building with gcc didn't work if gcc had been configured to use the native assembler, which is becoming more common.

- On AIX, test the `-qlonglong` compiler option rather than just assuming it's safe to use (Noah Misch)
- On AIX, use `-Wl,-brtllib` link option to allow symbols to be resolved at runtime (Noah Misch)

Perl relies on this ability in 5.8.0 and later.

- Avoid use of inline functions when compiling with 32-bit xlc, due to compiler bugs (Noah Misch)
- Use `librt` for `sched_yield()` when necessary, which it is on some Solaris versions (Oskari Saarenmaa)
- Fix Windows `install.bat` script to handle target directory names that contain spaces (Heikki Linnakangas)
- Make the numeric form of the PostgreSQL version number (e.g., `90405`) readily available to extension Makefiles, as a variable named `VERSION_NUM` (Michael Paquier)
- Update time zone data files to tzdata release 2015g for DST law changes in Cayman Islands, Fiji, Moldova, Morocco, Norfolk Island, North Korea, Turkey, and Uruguay. There is a new zone name `America/Fort_Nelson` for the Canadian Northern Rockies.

E.136. Release 9.0.22

Release date: 2015-06-12

This release contains a small number of fixes from 9.0.21. For information about new features in the 9.0 major release, see Section E.158.

The PostgreSQL community will stop releasing updates for the 9.0.X release series in September 2015. Users are encouraged to update to a newer release branch soon.

E.136.1. Migration to Version 9.0.22

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.18, see Section E.140.

E.136.2. Changes

- Fix rare failure to invalidate relation cache init file (Tom Lane)

  With just the wrong timing of concurrent activity, a `VACUUM FULL` on a system catalog might fail to update the “init file” that's used to avoid cache-loading work for new sessions. This would result in later sessions being unable to access that catalog at all. This is a very ancient bug, but it's so hard to trigger that no reproducible case had been seen until recently.

- Avoid deadlock between incoming sessions and `CREATE/DROP DATABASE` (Tom Lane)

  A new session starting in a database that is the target of a `DROP DATABASE` command, or is the template for a `CREATE DATABASE` command, could cause the command to wait for five seconds and then fail, even if the new session would have exited before that.

E.137. Release 9.0.21

Release date: 2015-06-04
This release contains a small number of fixes from 9.0.20. For information about new features in the 9.0 major release, see Section E.158.

The PostgreSQL community will stop releasing updates for the 9.0.X release series in September 2015. Users are encouraged to update to a newer release branch soon.

### E.137.1. Migration to Version 9.0.21

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.18, see Section E.140.

### E.137.2. Changes

- **Avoid failures while fsync'ing data directory during crash restart (Abhijit Menon-Sen, Tom Lane)**
  
  In the previous minor releases we added a patch to fsync everything in the data directory after a crash. Unfortunately its response to any error condition was to fail, thereby preventing the server from starting up, even when the problem was quite harmless. An example is that an unwritable file in the data directory would prevent restart on some platforms; but it is common to make SSL certificate files unwritable by the server. Revise this behavior so that permissions failures are ignored altogether, and other types of failures are logged but do not prevent continuing.

- **Remove configure's check prohibiting linking to a threaded libpython on OpenBSD (Tom Lane)**
  
  The failure this restriction was meant to prevent seems to not be a problem anymore on current OpenBSD versions.

- **Allow libpq to use TLS protocol versions beyond v1 (Noah Misch)**
  
  For a long time, libpq was coded so that the only SSL protocol it would allow was TLS v1. Now that newer TLS versions are becoming popular, allow it to negotiate the highest commonly-supported TLS version with the server. (PostgreSQL servers were already capable of such negotiation, so no change is needed on the server side.) This is a back-patch of a change already released in 9.4.0.

### E.138. Release 9.0.20

**Release date:** 2015-05-22

This release contains a variety of fixes from 9.0.19. For information about new features in the 9.0 major release, see Section E.158.

The PostgreSQL community will stop releasing updates for the 9.0.X release series in September 2015. Users are encouraged to update to a newer release branch soon.

### E.138.1. Migration to Version 9.0.20

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.18, see Section E.140.

### E.138.2. Changes

- **Avoid possible crash when client disconnects just before the authentication timeout expires (Benkocs Norbert Attila)**
  
  If the timeout interrupt fired partway through the session shutdown sequence, SSL-related state would be freed twice, typically causing a crash and hence denial of service to other sessions. Ex-
 experimentation shows that an unauthenticated remote attacker could trigger the bug somewhat consistently, hence treat as security issue. (CVE-2015-3165)

• Improve detection of system-call failures (Noah Misch)

Our replacement implementation of `snprintf()` failed to check for errors reported by the underlying system library calls; the main case that might be missed is out-of-memory situations. In the worst case this might lead to information exposure, due to our code assuming that a buffer had been overwritten when it hadn't been. Also, there were a few places in which security-relevant calls of other system library functions did not check for failure.

It remains possible that some calls of the `*printf()` family of functions are vulnerable to information disclosure if an out-of-memory error occurs at just the wrong time. We judge the risk to not be large, but will continue analysis in this area. (CVE-2015-3166)

• In `contrib/pgcrypto`, uniformly report decryption failures as “Wrong key or corrupt data” (Noah Misch)

Previously, some cases of decryption with an incorrect key could report other error message texts. It has been shown that such variance in error reports can aid attackers in recovering keys from other systems. While it's unknown whether `pgcrypto`'s specific behaviors are likewise exploitable, it seems better to avoid the risk by using a one-size-fits-all message. (CVE-2015-3167)

• Fix incorrect checking of deferred exclusion constraints after a HOT update (Tom Lane)

If a new row that potentially violates a deferred exclusion constraint is HOT-updated (that is, no indexed columns change and the row can be stored back onto the same table page) later in the same transaction, the exclusion constraint would be reported as violated when the check finally occurred, even if the row(s) the new row originally conflicted with had been deleted.

• Prevent improper reordering of antijoins (NOT EXISTS joins) versus other outer joins (Tom Lane)

This oversight in the planner has been observed to cause “could not find RelOptInfo for given relids” errors, but it seems possible that sometimes an incorrect query plan might get past that consistency check and result in silently-wrong query output.

• Fix incorrect matching of subexpressions in outer-join plan nodes (Tom Lane)

Previously, if textually identical non-strict subexpressions were used both above and below an outer join, the planner might try to re-use the value computed below the join, which would be incorrect because the executor would force the value to NULL in case of an unmatched outer row.

• Fix GEQO planner to cope with failure of its join order heuristic (Tom Lane)

This oversight has been seen to lead to “failed to join all relations together” errors in queries involving `LATERAL`, and that might happen in other cases as well.

• Fix possible deadlock at startup when `max_prepared_transactions` is too small (Heikki Linnakangas)

• Don't archive useless preallocated WAL files after a timeline switch (Heikki Linnakangas)

• Avoid “cannot GetMultiXactIdMembers() during recovery” error (Álvaro Herrera)

• Recursively `fsync()` the data directory after a crash (Abhijit Menon-Sen, Robert Haas)

This ensures consistency if another crash occurs shortly later. (The second crash would have to be a system-level crash, not just a database crash, for there to be a problem.)

• Fix autovacuum launcher's possible failure to shut down, if an error occurs after it receives `SIGTERM` (Álvaro Herrera)
• Cope with unexpected signals in LockBufferForCleanup() (Andres Freund)

This oversight could result in spurious errors about “multiple backends attempting to wait for pin-count 1”.

• Avoid waiting for WAL flush or synchronous replication during commit of a transaction that was read-only so far as the user is concerned (Andres Freund)

Previously, a delay could occur at commit in transactions that had written WAL due to HOT page pruning, leading to undesirable effects such as sessions getting stuck at startup if all synchronous replicas are down. Sessions have also been observed to get stuck in catchup interrupt processing when using synchronous replication; this will fix that problem as well.

• Fix crash when manipulating hash indexes on temporary tables (Heikki Linnakangas)

• Fix possible failure during hash index bucket split, if other processes are modifying the index concurrently (Tom Lane)

• Check for interrupts while analyzing index expressions (Jeff Janes)

   ANALYZE executes index expressions many times; if there are slow functions in such an expression, it's desirable to be able to cancel the ANALYZE before that loop finishes.

• Add the name of the target server to object description strings for foreign-server user mappings (Alvaro Herrera)

• Recommend setting include_realm to 1 when using Kerberos/GSSAPI/SSPI authentication (Stephen Frost)

   Without this, identically-named users from different realms cannot be distinguished. For the moment this is only a documentation change, but it will become the default setting in PostgreSQL 9.5.

• Remove code for matching IPv4 pg_hba.conf entries to IPv4-in-IPv6 addresses (Tom Lane)

   This hack was added in 2003 in response to a report that some Linux kernels of the time would report IPv4 connections as having IPv4-in-IPv6 addresses. However, the logic was accidentally broken in 9.0. The lack of any field complaints since then shows that it's not needed anymore. Now we have reports that the broken code causes crashes on some systems, so let's just remove it rather than fix it. (Had we chosen to fix it, that would make for a subtle and potentially security-sensitive change in the effective meaning of IPv4 pg_hba.conf entries, which does not seem like a good thing to do in minor releases.)

• While shutting down service on Windows, periodically send status updates to the Service Control Manager to prevent it from killing the service too soon; and ensure that pg_ctl will wait for shutdown (Krystian Bigaj)

• Reduce risk of network deadlock when using libpq's non-blocking mode (Heikki Linnakangas)

   When sending large volumes of data, it's important to drain the input buffer every so often, in case the server has sent enough response data to cause it to block on output. (A typical scenario is that the server is sending a stream of NOTICE messages during COPY FROM STDIN.) This worked properly in the normal blocking mode, but not so much in non-blocking mode. We've modified libpq to opportunistically drain input when it can, but a full defense against this problem requires application cooperation: the application should watch for socket read-ready as well as write-ready conditions, and be sure to call PQconsumeInput() upon read-ready.

• Fix array handling in ecpg (Michael Meskes)

• Fix psql to sanely handle URIs and conninfo strings as the first parameter to \connect (David Fetter, Andrew Dunstan, Alvaro Herrera)
This syntax has been accepted (but undocumented) for a long time, but previously some parameters might be taken from the old connection instead of the given string, which was agreed to be undesirable.

- Suppress incorrect complaints from psql on some platforms that it failed to write ~/.psql_history at exit (Tom Lane)
  
  This misbehavior was caused by a workaround for a bug in very old (pre-2006) versions of libedit. We fixed it by removing the workaround, which will cause a similar failure to appear for anyone still using such versions of libedit. Recommendation: upgrade that library, or use libreadline.

- Fix pg_dump's rule for deciding which casts are system-provided casts that should not be dumped (Tom Lane)

- Fix dumping of views that are just VALUES(...) but have column aliases (Tom Lane)

- In pg_upgrade, force timeline 1 in the new cluster (Bruce Momjian)
  
  This change prevents upgrade failures caused by bogus complaints about missing WAL history files.

- In pg_upgrade, check for improperly non-connectable databases before proceeding (Bruce Momjian)

- In pg_upgrade, quote directory paths properly in the generated delete_old_cluster script (Bruce Momjian)

- In pg_upgrade, preserve database-level freezing info properly (Bruce Momjian)
  
  This oversight could cause missing-clog-file errors for tables within the postgres and template databases.

- Run pg_upgrade and pg_resetxlog with restricted privileges on Windows, so that they don't fail when run by an administrator (Muhammad Asif Naeem)

- Fix slow sorting algorithm in contrib/intarray (Tom Lane)

- Fix compile failure on Sparc V8 machines (Rob Rowan)

- Update time zone data files to tzdata release 2015d for DST law changes in Egypt, Mongolia, and Palestine, plus historical changes in Canada and Chile. Also adopt revised zone abbreviations for the America/Adak zone (HST/HDT not HAST/HADT).

E.139. Release 9.0.19

Release date: 2015-02-05

This release contains a variety of fixes from 9.0.18. For information about new features in the 9.0 major release, see Section E.158.

E.139.1. Migration to Version 9.0.19

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.18, see Section E.140.

E.139.2. Changes

- Fix buffer overruns in to_char() (Bruce Momjian)
When `to_char()` processes a numeric formatting template calling for a large number of digits, PostgreSQL would read past the end of a buffer. When processing a crafted timestamp formatting template, PostgreSQL would write past the end of a buffer. Either case could crash the server. We have not ruled out the possibility of attacks that lead to privilege escalation, though they seem unlikely. (CVE-2015-0241)

• Fix buffer overrun in replacement `printf()` functions (Tom Lane)

PostgreSQL includes a replacement implementation of `printf` and related functions. This code will overrun a stack buffer when formatting a floating point number (conversion specifiers e, E, f, F, g or G) with requested precision greater than about 500. This will crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. A database user can trigger such a buffer overrun through the `to_char()` SQL function. While that is the only affected core PostgreSQL functionality, extension modules that use printf-family functions may be at risk as well.

This issue primarily affects PostgreSQL on Windows. PostgreSQL uses the system implementation of these functions where adequate, which it is on other modern platforms. (CVE-2015-0242)

• Fix buffer overruns in `contrib/pgcrypto` (Marko Tiikkaja, Noah Misch)

Errors in memory size tracking within the `pgcrypto` module permitted stack buffer overruns and improper dependence on the contents of uninitialized memory. The buffer overrun cases can crash the server, and we have not ruled out the possibility of attacks that lead to privilege escalation. (CVE-2015-0243)

• Fix possible loss of frontend/backend protocol synchronization after an error (Heikki Linnakangas)

If any error occurred while the server was in the middle of reading a protocol message from the client, it could lose synchronization and incorrectly try to interpret part of the message's data as a new protocol message. An attacker able to submit crafted binary data within a command parameter might succeed in injecting his own SQL commands this way. Statement timeout and query cancellation are the most likely sources of errors triggering this scenario. Particularly vulnerable are applications that use a timeout and also submit arbitrary user-crafted data as binary query parameters. Disabling statement timeout will reduce, but not eliminate, the risk of exploit. Our thanks to Emil Lenngren for reporting this issue. (CVE-2015-0244)

• Fix information leak via constraint-violation error messages (Stephen Frost)

Some server error messages show the values of columns that violate a constraint, such as a unique constraint. If the user does not have `SELECT` privilege on all columns of the table, this could mean exposing values that the user should not be able to see. Adjust the code so that values are displayed only when they came from the SQL command or could be selected by the user. (CVE-2014-8161)

• Lock down regression testing's temporary installations on Windows (Noah Misch)

Use SSPI authentication to allow connections only from the OS user who launched the test suite. This closes on Windows the same vulnerability previously closed on other platforms, namely that other users might be able to connect to the test postmaster. (CVE-2014-0067)

• Avoid possible data corruption if `ALTER DATABASE SET TABLESPACE` is used to move a database to a new tablespace and then shortly later move it back to its original tablespace (Tom Lane)

• Avoid corrupting tables when `ANALYZE` inside a transaction is rolled back (Andres Freund, Tom Lane, Michael Paquier)

If the failing transaction had earlier removed the last index, rule, or trigger from the table, the table would be left in a corrupted state with the relevant `pg_class` flags not set though they should be.

• Fix use-of-already-freed-memory problem in EvalPlanQual processing (Tom Lane)
In **READ COMMITTED** mode, queries that lock or update recently-updated rows could crash as a result of this bug.

- Fix planning of **SELECT FOR UPDATE** when using a partial index on a child table (Kyotaro Horiguchi)

  In **READ COMMITTED** mode, **SELECT FOR UPDATE** must also recheck the partial index’s **WHERE** condition when rechecking a recently-updated row to see if it still satisfies the query’s **WHERE** condition. This requirement was missed if the index belonged to an inheritance child table, so that it was possible to incorrectly return rows that no longer satisfy the query condition.

- Fix corner case wherein **SELECT FOR UPDATE** could return a row twice, and possibly miss returning other rows (Tom Lane)

  In **READ COMMITTED** mode, a **SELECT FOR UPDATE** that is scanning an inheritance tree could incorrectly return a row from a prior child table instead of the one it should return from a later child table.

- Reject duplicate column names in the referenced-columns list of a **FOREIGN KEY** declaration (David Rowley)

  This restriction is per SQL standard. Previously we did not reject the case explicitly, but later on the code would fail with bizarre-looking errors.

- Fix bugs in raising a numeric value to a large integral power (Tom Lane)

  The previous code could get a wrong answer, or consume excessive amounts of time and memory before realizing that the answer must overflow.

- In **numeric_recv()**, truncate away any fractional digits that would be hidden according to the value’s **dscale** field (Tom Lane)

  A numeric value’s display scale (**dscale**) should never be less than the number of nonzero fractional digits; but apparently there’s at least one broken client application that transmits binary numeric values in which that’s true. This leads to strange behavior since the extra digits are taken into account by arithmetic operations even though they aren’t printed. The least risky fix seems to be to truncate away such “hidden” digits on receipt, so that the value is indeed what it prints as.

- Reject out-of-range numeric timezone specifications (Tom Lane)

  Simple numeric timezone specifications exceeding +/- 168 hours (one week) would be accepted, but could then cause null-pointer dereference crashes in certain operations. There’s no use-case for such large UTC offsets, so reject them.

- Fix bugs in **tsquery @> tsquery** operator (Heikki Linnakangas)

  Two different terms would be considered to match if they had the same CRC. Also, if the second operand had more terms than the first, it would be assumed not to be contained in the first; which is wrong since it might contain duplicate terms.

- Improve ispell dictionary’s defenses against bad affix files (Tom Lane)

- Allow more than 64K phrases in a thesaurus dictionary (David Boutin)

  The previous coding could crash on an oversize dictionary, so this was deemed a back-patchable bug fix rather than a feature addition.

- Fix namespace handling in **xpath()** (Ali Akbar)

  Previously, the **xml** value resulting from an **xpath()** call would not have namespace declarations if the namespace declarations were attached to an ancestor element in the input **xml** value, rather
than to the specific element being returned. Propagate the ancestral declaration so that the result is
correct when considered in isolation.

• Fix planner problems with nested append relations, such as inherited tables within \texttt{UNION ALL}
subqueries (Tom Lane)

• Fail cleanly when a GiST index tuple doesn't fit on a page, rather than going into infinite recursion
(Andrew Gierth)

• Exempt tables that have per-table \texttt{cost\_limit} and/or \texttt{cost\_delay} settings from autovacuum's
  global cost balancing rules (Álvaro Herrera)

The previous behavior resulted in basically ignoring these per-table settings, which was unintended.
Now, a table having such settings will be vacuumed using those settings, independently of what
is going on in other autovacuum workers. This may result in heavier total I/O load than before, so
such settings should be re-examined for sanity.

• Avoid wholesale autovacuuming when autovacuum is nominally off (Tom Lane)

Even when autovacuum is nominally off, we will still launch autovacuum worker processes to
vacuum tables that are at risk of XID wraparound. However, such a worker process then proceeded
to vacuum all tables in the target database, if they met the usual thresholds for autovacuuming. This
is at best pretty unexpected; at worst it delays response to the wraparound threat. Fix it so that if
autovacuum is turned off, workers only do anti-wraparound vacuums and not any other work.

• Fix race condition between hot standby queries and replaying a full-page image (Heikki Linnakan-
gas)

This mistake could result in transient errors in queries being executed in hot standby.

• Fix several cases where recovery logic improperly ignored WAL records for \texttt{COMMIT/ABORT PREPARED}
  (Heikki Linnakangas)

The most notable oversight was that \texttt{recovery\_target\_xid} could not be used to stop at a two-
phase commit.

• Avoid creating unnecessary \texttt{.ready} marker files for timeline history files (Fujii Masao)

• Fix possible null pointer dereference when an empty prepared statement is used and the \texttt{log\_s-
tatement} setting is \texttt{mod} or \texttt{ddl} (Fujii Masao)

• Change “pgstat wait timeout” warning message to be \texttt{LOG} level, and rephrase it to be more under-
  standable (Tom Lane)

This message was originally thought to be essentially a can't-happen case, but it occurs often enough
on our slower buildfarm members to be a nuisance. Reduce it to \texttt{LOG} level, and expend a bit more
effort on the wording: it now reads “using stale statistics instead of current ones because stats col-
lector is not responding”.

• Fix SPARC spinlock implementation to ensure correctness if the CPU is being run in a non-TSO
  coherency mode, as some non-Solaris kernels do (Andres Freund)

• Warn if macOS's \texttt{setlocale()} starts an unwanted extra thread inside the postmaster (Noah
  Misch)

• Fix processing of repeated \texttt{dbname} parameters in \texttt{PQconnectdbParams()} (Alex Shulgin)

Unexpected behavior ensued if the first occurrence of \texttt{dbname} contained a connection string or
URI to be expanded.

• Ensure that \texttt{libpq} reports a suitable error message on unexpected socket EOF (Marko Tiikkaja, Tom
  Lane)
Depending on kernel behavior, libpq might return an empty error string rather than something useful when the server unexpectedly closed the socket.

- Clear any old error message during PQreset() (Heikki Linnakangas)

  If PQreset() is called repeatedly, and the connection cannot be re-established, error messages from the failed connection attempts kept accumulating in the PGconn's error string.

- Properly handle out-of-memory conditions while parsing connection options in libpq (Alex Shulgin, Heikki Linnakangas)

- Fix array overrun in ecpg's version of ParseDateTime() (Michael Paquier)

- In initdb, give a clearer error message if a password file is specified but is empty (Mats Erik Andersson)

- Fix psql's \s command to work nicely with libedit, and add pager support (Stepan Rutz, Tom Lane)

  When using libedit rather than readline, \s printed the command history in a fairly unreadable encoded format, and on recent libedit versions might fail altogether. Fix that by printing the history ourselves rather than having the library do it. A pleasant side-effect is that the pager is used if appropriate.

  This patch also fixes a bug that caused newline encoding to be applied inconsistently when saving the command history with libedit. Multiline history entries written by older psql versions will be read cleanly with this patch, but perhaps not vice versa, depending on the exact libedit versions involved.

- Improve consistency of parsing of psql's special variables (Tom Lane)

  Allow variant spellings of on and off (such as 1/0) for ECHO_HIDDEN and ON_ERROR_ROLLBACK. Report a warning for unrecognized values for COMP_KEYWORD_CASE, ECHO, ECHO_HIDDEN, HISTCONTROL, ON_ERROR_ROLLBACK, and VERBOSITY. Recognize all values for all these variables case-insensitively; previously there was a mishmash of case-sensitive and case-insensitive behaviors.

- Fix psql's expanded-mode display to work consistently when using border = 3 and linestyle = ascii or unicode (Stephen Frost)

- Fix possible deadlock during parallel restore of a schema-only dump (Robert Haas, Tom Lane)

- Fix core dump in pg_dump --binary-upgrade on zero-column composite type (Rushabh Lathia)

- Fix block number checking in contrib/pageinspect's get_raw_page() (Tom Lane)

  The incorrect checking logic could prevent access to some pages in non-main relation forks.

- Fix contrib/pgcrypto's pgp_sym_decrypt() to not fail on messages whose length is 6 less than a power of 2 (Marko Tiikkaja)

- Handle unexpected query results, especially NULLs, safely in contrib/tablefunc's connectby() (Michael Paquier)

  connectby() previously crashed if it encountered a NULL key value. It now prints that row but doesn't recurse further.

- Avoid a possible crash in contrib/xml2's xslt_process() (Mark Simonetti)

  libxml seems to have an undocumented dependency on the order in which resources are freed; reorder our calls to avoid a crash.
• Numerous cleanups of warnings from Coverity static code analyzer (Andres Freund, Tatsuo Ishii, Marko Kreen, Tom Lane, Michael Paquier)

These changes are mostly cosmetic but in some cases fix corner-case bugs, for example a crash rather than a proper error report after an out-of-memory failure. None are believed to represent security issues.

• Detect incompatible OpenLDAP versions during build (Noah Misch)

With OpenLDAP versions 2.4.24 through 2.4.31, inclusive, PostgreSQL backends can crash at exit. Raise a warning during configure based on the compile-time OpenLDAP version number, and test the crashing scenario in the contrib/dblink regression test.

• In non-MSVC Windows builds, ensure libpq.dll is installed with execute permissions (Noah Misch)

• Make pg_regress remove any temporary installation it created upon successful exit (Tom Lane)

This results in a very substantial reduction in disk space usage during make check-world, since that sequence involves creation of numerous temporary installations.

• Support time zone abbreviations that change UTC offset from time to time (Tom Lane)

Previously, PostgreSQL assumed that the UTC offset associated with a time zone abbreviation (such as EST) never changes in the usage of any particular locale. However this assumption fails in the real world, so introduce the ability for a zone abbreviation to represent a UTC offset that sometimes changes. Update the zone abbreviation definition files to make use of this feature in timezone locales that have changed the UTC offset of their abbreviations since 1970 (according to the IANA timezone database). In such timezones, PostgreSQL will now associate the correct UTC offset with the abbreviation depending on the given date.

• Update time zone abbreviations lists (Tom Lane)

Add CST (China Standard Time) to our lists. Remove references to ADT as “Arabia Daylight Time”, an abbreviation that's been out of use since 2007; therefore, claiming there is a conflict with “Atlantic Daylight Time” doesn't seem especially helpful. Fix entirely incorrect GMT offsets for CKT (Cook Islands), FJT, and FJST (Fiji); we didn't even have them on the proper side of the date line.

• Update time zone data files to tzdata release 2015a.

The IANA timezone database has adopted abbreviations of the form A\times7/A\timesDT for all Australian time zones, reflecting what they believe to be current majority practice Down Under. These names do not conflict with usage elsewhere (other than ACST for Acre Summer Time, which has been in disuse since 1994). Accordingly, adopt these names into our “Default” timezone abbreviation set. The “Australia” abbreviation set now contains only CST, EAST, EST, SAST, SAT, and WST, all of which are thought to be mostly historical usage. Note that SAST has also been changed to be South Africa Standard Time in the “Default” abbreviation set.

Also, add zone abbreviations SRET (Asia/Srednekolymsk) and XJT (Asia/Urumqi), and use WSST/ WSDT for western Samoa. Also, there were DST law changes in Chile, Mexico, the Turks & Caicos Islands (America/Grand_Turk), and Fiji. There is a new zone Pacific/Bougainville for portions of Papua New Guinea. Also, numerous corrections for historical (pre-1970) time zone data.

**E.140. Release 9.0.18**

**Release date:** 2014-07-24

This release contains a variety of fixes from 9.0.17. For information about new features in the 9.0 major release, see Section E.158.
E.140.1. Migration to Version 9.0.18

A dump/restore is not required for those running 9.0.X.

However, this release corrects an index corruption problem in some GiST indexes. See the first changelog entry below to find out whether your installation has been affected and what steps you should take if so.

Also, if you are upgrading from a version earlier than 9.0.15, see Section E.143.

E.140.2. Changes

- Correctly initialize padding bytes in contrib/btree_gist indexes on bit columns (Heikki Linnakangas)
  
  This error could result in incorrect query results due to values that should compare equal not being seen as equal. Users with GiST indexes on bit or bit varying columns should REINDEX those indexes after installing this update.

- Protect against torn pages when deleting GIN list pages (Heikki Linnakangas)
  
  This fix prevents possible index corruption if a system crash occurs while the page update is being written to disk.

- Don’t clear the right-link of a GiST index page while replaying updates from WAL (Heikki Linnakangas)
  
  This error could lead to transiently wrong answers from GiST index scans performed in Hot Standby.

- Fix possibly-incorrect cache invalidation during nested calls to ReceiveSharedInvalidMessages (Andres Freund)

- Don’t assume a subquery’s output is unique if there’s a set-returning function in its targetlist (David Rowley)
  
  This oversight could lead to misoptimization of constructs like WHERE x IN (SELECT y, generate_series(1,10) FROM t GROUP BY y).

- Fix failure to detoast fields in composite elements of structured types (Tom Lane)
  
  This corrects cases where TOAST pointers could be copied into other tables without being dereferenced. If the original data is later deleted, it would lead to errors like “missing chunk number 0 for toast value ...” when the now-dangling pointer is used.

- Fix “record type has not been registered” failures with whole-row references to the output of Append plan nodes (Tom Lane)

- Fix possible crash when invoking a user-defined function while rewinding a cursor (Tom Lane)

- Fix query-lifespan memory leak while evaluating the arguments for a function in FROM (Tom Lane)

- Fix session-lifespan memory leaks in regular-expression processing (Tom Lane, Arthur O’Dwyer, Greg Stark)

- Fix data encoding error in hungarian.stop (Tom Lane)

- Fix liveness checks for rows that were inserted in the current transaction and then deleted by a now-rolled-back subtransaction (Andres Freund)
This could cause problems (at least spurious warnings, and at worst an infinite loop) if CREATE INDEX or CLUSTER were done later in the same transaction.

- Clear pg_stat_activity.xact_start during PREPARE TRANSACTION (Andres Freund)
  After the PREPARE, the originating session is no longer in a transaction, so it should not continue to display a transaction start time.

- Fix REASSIGN OWNED to not fail for text search objects (Álvaro Herrera)

- Block signals during postmaster startup (Tom Lane)
  This ensures that the postmaster will properly clean up after itself if, for example, it receives SIGINT while still starting up.

- Secure Unix-domain sockets of temporary postmasters started during make check (Noah Misch)
  Any local user able to access the socket file could connect as the server's bootstrap superuser, then proceed to execute arbitrary code as the operating-system user running the test, as we previously noted in CVE-2014-0067. This change defends against that risk by placing the server's socket in a temporary, mode 0700 subdirectory of /tmp. The hazard remains however on platforms where Unix sockets are not supported, notably Windows, because then the temporary postmaster must accept local TCP connections.

  A useful side effect of this change is to simplify make check testing in builds that override DEFAULT_PGSOCKET_DIR. Popular non-default values like /var/run/postgresql are often not writable by the build user, requiring workarounds that will no longer be necessary.

- Fix tablespace creation WAL replay to work on Windows (MauMau)

- Fix detection of socket creation failures on Windows (Bruce Momjian)

- On Windows, allow new sessions to absorb values of PGC_BACKEND parameters (such as log_connections) from the configuration file (Amit Kapila)
  Previously, if such a parameter were changed in the file post-startup, the change would have no effect.

- Properly quote executable path names on Windows (Nikhil Deshpande)
  This oversight could cause initdb and pg_upgrade to fail on Windows, if the installation path contained both spaces and @ signs.

- Fix linking of libpython on macOS (Tom Lane)
  The method we previously used can fail with the Python library supplied by Xcode 5.0 and later.

- Avoid buffer bloat in libpq when the server consistently sends data faster than the client can absorb it (Shin-ichi Morita, Tom Lane)
  Libpq could be coerced into enlarging its input buffer until it runs out of memory (which would be reported misleadingly as “lost synchronization with server”). Under ordinary circumstances it's quite far-fetched that data could be continuously transmitted more quickly than the recv() loop can absorb it, but this has been observed when the client is artificially slowed by scheduler constraints.

- Ensure that LDAP lookup attempts in libpq time out as intended (Laurenz Albe)

- Fix ecpg to do the right thing when an array of char * is the target for a FETCH statement returning more than one row, as well as some other array-handling fixes (Ashutosh Bapat)

- Fix pg_restore's processing of old-style large object comments (Tom Lane)
A direct-to-database restore from an archive file generated by a pre-9.0 version of pg_dump would usually fail if the archive contained more than a few comments for large objects.

- In contrib/pgcrypto functions, ensure sensitive information is cleared from stack variables before returning (Marko Kreen)
- In contrib/uuid-ossp, cache the state of the OSSP UUID library across calls (Tom Lane)
  This improves the efficiency of UUID generation and reduces the amount of entropy drawn from /dev/urandom, on platforms that have that.
- Update time zone data files to tzdata release 2014e for DST law changes in Crimea, Egypt, and Morocco.

E.141. Release 9.0.17

Release date: 2014-03-20

This release contains a variety of fixes from 9.0.16. For information about new features in the 9.0 major release, see Section E.158.

E.141.1. Migration to Version 9.0.17

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.15, see Section E.143.

E.141.2. Changes

- Restore GIN metapages unconditionally to avoid torn-page risk (Heikki Linnakangas)
  Although this oversight could theoretically result in a corrupted index, it is unlikely to have caused any problems in practice, since the active part of a GIN metapage is smaller than a standard 512-byte disk sector.
- Avoid race condition in checking transaction commit status during receipt of a NOTIFY message (Marko Tiikkaja)
  This prevents a scenario wherein a sufficiently fast client might respond to a notification before database updates made by the notifier have become visible to the recipient.
- Allow regular-expression operators to be terminated early by query cancel requests (Tom Lane)
  This prevents scenarios wherein a pathological regular expression could lock up a server process uninterruptibly for a long time.
- Remove incorrect code that tried to allow OVERLAPS with single-element row arguments (Joshua Yanovski)
  This code never worked correctly, and since the case is neither specified by the SQL standard nor documented, it seemed better to remove it than fix it.
- Avoid getting more than AccessShareLock when de-parsing a rule or view (Dean Rasheed)
  This oversight resulted in pg_dump unexpectedly acquiring RowExclusiveLock locks on tables mentioned as the targets of INSERT/UPDATE/DELETE commands in rules. While usually harmless, that could interfere with concurrent transactions that tried to acquire, for example, ShareLock on those tables.
- Improve performance of index endpoint probes during planning (Tom Lane)
This change fixes a significant performance problem that occurred when there were many not-yet-committed rows at the end of the index, which is a common situation for indexes on sequentially-assigned values such as timestamps or sequence-generated identifiers.

- Fix test to see if hot standby connections can be allowed immediately after a crash (Heikki Linnakangas)
- Prevent interrupts while reporting non-ERROR messages (Tom Lane)
  This guards against rare server-process freezeups due to recursive entry to \texttt{syslog()}, and perhaps other related problems.
- Prevent intermittent “could not reserve shared memory region” failures on recent Windows versions (MauMau)
- Update time zone data files to \texttt{tzdata} release 2014a for DST law changes in Fiji and Turkey, plus historical changes in Israel and Ukraine.

**E.142. Release 9.0.16**

**Release date:** 2014-02-20

This release contains a variety of fixes from 9.0.15. For information about new features in the 9.0 major release, see Section E.158.

**E.142.1. Migration to Version 9.0.16**

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.15, see Section E.143.

**E.142.2. Changes**

- Shore up \texttt{GRANT ... WITH ADMIN OPTION} restrictions (Noah Misch)
  Granting a role without \texttt{ADMIN OPTION} is supposed to prevent the grantee from adding or removing members from the granted role, but this restriction was easily bypassed by doing \texttt{SET ROLE} first. The security impact is mostly that a role member can revoke the access of others, contrary to the wishes of his grantor. Unapproved role member additions are a lesser concern, since an uncooperative role member could provide most of his rights to others anyway by creating views or \texttt{SECURITY DEFINER} functions. (CVE-2014-0060)

- Prevent privilege escalation via manual calls to PL validator functions (Andres Freund)
  The primary role of PL validator functions is to be called implicitly during \texttt{CREATE FUNCTION}, but they are also normal SQL functions that a user can call explicitly. Calling a validator on a function actually written in some other language was not checked for and could be exploited for privilege-escalation purposes. The fix involves adding a call to a privilege-checking function in each validator function. Non-core procedural languages will also need to make this change to their own validator functions, if any. (CVE-2014-0061)

- Avoid multiple name lookups during table and index DDL (Robert Haas, Andres Freund)
  If the name lookups come to different conclusions due to concurrent activity, we might perform some parts of the DDL on a different table than other parts. At least in the case of \texttt{CREATE INDEX}, this can be used to cause the permissions checks to be performed against a different table than the index creation, allowing for a privilege escalation attack. (CVE-2014-0062)

- Prevent buffer overrun with long datetime strings (Noah Misch)
Release Notes

The `MAXDATELEN` constant was too small for the longest possible value of type `interval`, allowing a buffer overrun in `interval_out()`. Although the datetime input functions were more careful about avoiding buffer overrun, the limit was short enough to cause them to reject some valid inputs, such as input containing a very long timezone name. The `ecpg` library contained these vulnerabilities along with some of its own. (CVE-2014-0063)

- Prevent buffer overrun due to integer overflow in size calculations (Noah Misch, Heikki Linnakangas)

Several functions, mostly type input functions, calculated an allocation size without checking for overflow. If overflow did occur, a too-small buffer would be allocated and then written past. (CVE-2014-0064)

- Prevent overruns of fixed-size buffers (Peter Eisentraut, Jozef Mlich)

Use `strlcpy()` and related functions to provide a clear guarantee that fixed-size buffers are not overrun. Unlike the preceding items, it is unclear whether these cases really represent live issues, since in most cases there appear to be previous constraints on the size of the input string. Nonetheless it seems prudent to silence all Coverity warnings of this type. (CVE-2014-0065)

- Avoid crashing if `crypt()` returns NULL (Honza Horak, Bruce Momjian)

There are relatively few scenarios in which `crypt()` could return NULL, but `contrib/chkpass` would crash if it did. One practical case in which this could be an issue is if libc is configured to refuse to execute unapproved hashing algorithms (e.g., “FIPS mode”). (CVE-2014-0066)

- Document risks of `make check` in the regression testing instructions (Noah Misch, Tom Lane)

Since the temporary server started by `make check` uses “trust” authentication, another user on the same machine could connect to it as database superuser, and then potentially exploit the privileges of the operating-system user who started the tests. A future release will probably incorporate changes in the testing procedure to prevent this risk, but some public discussion is needed first. So for the moment, just warn people against using `make check` when there are untrusted users on the same machine. (CVE-2014-0067)

- Fix possible mis-replay of WAL records when some segments of a relation aren't full size (Greg Stark, Tom Lane)

The WAL update could be applied to the wrong page, potentially many pages past where it should have been. Aside from corrupting data, this error has been observed to result in significant “bloat” of standby servers compared to their masters, due to updates being applied far beyond where the end-of-file should have been. This failure mode does not appear to be a significant risk during crash recovery, only when initially synchronizing a standby created from a base backup taken from a quickly-changing master.

- Fix bug in determining when recovery has reached consistency (Tomonari Katsumata, Heikki Linnakangas)

In some cases WAL replay would mistakenly conclude that the database was already consistent at the start of replay, thus possibly allowing hot-standby queries before the database was really consistent. Other symptoms such as “PANIC: WAL contains references to invalid pages” were also possible.

- Fix improper locking of btree index pages while replaying a `VACUUM` operation in hot-standby mode (Andres Freund, Heikki Linnakangas, Tom Lane)

This error could result in “PANIC: WAL contains references to invalid pages” failures.

- Ensure that insertions into non-leaf GIN index pages write a full-page WAL record when appropriate (Heikki Linnakangas)
The previous coding risked index corruption in the event of a partial-page write during a system crash.

- Fix race conditions during server process exit (Robert Haas)

  Ensure that signal handlers don’t attempt to use the process’s MyProc pointer after it’s no longer valid.

- Fix unsafe references to errno within error reporting logic (Christian Kruse)

  This would typically lead to odd behaviors such as missing or inappropriate HINT fields.

- Fix possible crashes from using ereport() too early during server startup (Tom Lane)

  The principal case we’ve seen in the field is a crash if the server is started in a directory it doesn’t have permission to read.

- Clear retry flags properly in OpenSSL socket write function (Alexander Kukushkin)

  This omission could result in a server lockup after unexpected loss of an SSL-encrypted connection.

- Fix length checking for Unicode identifiers (U&"..." syntax) containing escapes (Tom Lane)

  A spurious truncation warning would be printed for such identifiers if the escaped form of the identifier was too long, but the identifier actually didn’t need truncation after de-escaping.

- Allow keywords that are type names to be used in lists of roles (Stephen Frost)

  A previous patch allowed such keywords to be used without quoting in places such as role identifiers; but it missed cases where a list of role identifiers was permitted, such as DROP ROLE.

- Fix possible crash due to invalid plan for nested sub-selects, such as WHERE (... x IN (SELECT ...) ...) IN (SELECT ...) (Tom Lane)

- Ensure that ANALYZE creates statistics for a table column even when all the values in it are “too wide” (Tom Lane)

  ANALYZE intentionally omits very wide values from its histogram and most-common-values calculations, but it neglected to do something sane in the case that all the sampled entries are too wide.

- In ALTER TABLE ... SET TABLESPACE, allow the database’s default tablespace to be used without a permissions check (Stephen Frost)

  CREATE TABLE has always allowed such usage, but ALTER TABLE didn’t get the memo.

- Fix “cannot accept a set” error when some arms of a CASE return a set and others don’t (Tom Lane)

- Fix checks for all-zero client addresses in pgstat functions (Kevin Grittner)

- Fix possible misclassification of multibyte characters by the text search parser (Tom Lane)

  Non-ASCII characters could be misclassified when using C locale with a multibyte encoding. On Cygwin, non-C locales could fail as well.

- Fix possible misbehavior in plainto_tsquery() (Heikki Linnakangas)

  Use memmove() not memcpy() for copying overlapping memory regions. There have been no field reports of this actually causing trouble, but it’s certainly risky.

- Accept SHIFT_JIS as an encoding name for locale checking purposes (Tatsuo Ishii)
Release Notes

- Fix misbehavior of PQhost() on Windows (Fujii Masao)

  It should return localhost if no host has been specified.

- Improve error handling in libpq and psql for failures during COPY TO STDOUT/FROM STDIN (Tom Lane)

  In particular this fixes an infinite loop that could occur in 9.2 and up if the server connection was lost during COPY FROM STDIN. Variants of that scenario might be possible in older versions, or with other client applications.

- Fix misaligned descriptors in ecpg (MauMau)

- In ecpg, handle lack of a hostname in the connection parameters properly (Michael Meskes)

- Fix performance regression in contrib/dblink connection startup (Joe Conway)

  Avoid an unnecessary round trip when client and server encodings match.

- In contrib/isn, fix incorrect calculation of the check digit for ISMN values (Fabien Coelho)

- Ensure client-code-only installation procedure works as documented (Peter Eisentraut)

- In Mingw and Cygwin builds, install the libpq DLL in the bin directory (Andrew Dunstan)

  This duplicates what the MSVC build has long done. It should fix problems with programs like psql failing to start because they can't find the DLL.

- Avoid using the deprecated dllwrap tool in Cygwin builds (Marco Atzeri)

- Don't generate plain-text HISTORY and src/test/regress/README files anymore (Tom Lane)

  These text files duplicated the main HTML and PDF documentation formats. The trouble involved in maintaining them greatly outweighs the likely audience for plain-text format. Distribution tarballs will still contain files by these names, but they'll just be stubs directing the reader to consult the main documentation. The plain-text INSTALL file will still be maintained, as there is arguably a use-case for that.

- Update time zone data files to tzdata release 2013i for DST law changes in Jordan and historical changes in Cuba.

  In addition, the zones Asia/Riyadh87, Asia/Riyadh88, and Asia/Riyadh89 have been removed, as they are no longer maintained by IANA, and never represented actual civil timekeeping practice.

E.143. Release 9.0.15

Release date: 2013-12-05

This release contains a variety of fixes from 9.0.14. For information about new features in the 9.0 major release, see Section E.158.

E.143.1. Migration to Version 9.0.15

A dump/restore is not required for those running 9.0.X.

However, this release corrects a number of potential data corruption issues. See the first two changelog entries below to find out whether your installation has been affected and what steps you can take if so.
Also, if you are upgrading from a version earlier than 9.0.13, see Section E.145.

### E.143.2. Changes

- **Fix VACUUM's tests to see whether it can update relfrozenxid (Andres Freund)**

  In some cases VACUUM (either manual or autovacuum) could incorrectly advance a table's relfrozenxid value, allowing tuples to escape freezing, causing those rows to become invisible once $2^{31}$ transactions have elapsed. The probability of data loss is fairly low since multiple incorrect advancements would need to happen before actual loss occurs, but it's not zero. Users upgrading from releases 9.0.4 or 8.4.8 or earlier are not affected, but all later versions contain the bug.

  The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix any latent corruption but will not be able to fix all pre-existing data errors. However, an installation can be presumed safe after performing this vacuuming if it has executed fewer than $2^{31}$ update transactions in its lifetime (check this with `SELECT txid_current() < 2^{31}`).

- **Fix initialization of pg_clog and pg_subtrans during hot standby startup (Andres Freund, Heikki Linnakangas)**

  This bug can cause data loss on standby servers at the moment they start to accept hot-standby queries, by marking committed transactions as uncommitted. The likelihood of such corruption is small unless, at the time of standby startup, the primary server has executed many updating transactions since its last checkpoint. Symptoms include missing rows, rows that should have been deleted being still visible, and obsolete versions of updated rows being still visible alongside their newer versions.

  This bug was introduced in versions 9.3.0, 9.2.5, 9.1.10, and 9.0.14. Standby servers that have only been running earlier releases are not at risk. It's recommended that standby servers that have ever run any of the buggy releases be re-cloned from the primary (e.g., with a new base backup) after upgrading.

- **Truncate pg_multixact contents during WAL replay (Andres Freund)**

  This avoids ever-increasing disk space consumption in standby servers.

- **Fix race condition in GIN index posting tree page deletion (Heikki Linnakangas)**

  This could lead to transient wrong answers or query failures.

- **Avoid flattening a subquery whose SELECT list contains a volatile function wrapped inside a sub-SELECT (Tom Lane)**

  This avoids unexpected results due to extra evaluations of the volatile function.

- **Fix planner's processing of non-simple-variable subquery outputs nested within outer joins (Tom Lane)**

  This error could lead to incorrect plans for queries involving multiple levels of subqueries within `JOIN` syntax.

- **Fix premature deletion of temporary files (Andres Freund)**

- **Fix possible read past end of memory in rule printing (Peter Eisentraut)**

- **Fix array slicing of `int2vector` and `oidvector` values (Tom Lane)**

  Expressions of this kind are now implicitly promoted to regular `int2` or `oid` arrays.

- **Fix incorrect behaviors when using a SQL-standard, simple GMT offset timezone (Tom Lane)**
In some cases, the system would use the simple GMT offset value when it should have used the regular timezone setting that had prevailed before the simple offset was selected. This change also causes the `timeofday` function to honor the simple GMT offset zone.

- Prevent possible misbehavior when logging translations of Windows error codes (Tom Lane)
- Properly quote generated command lines in `pg_ctl` (Naoya Anzai and Tom Lane)

This fix applies only to Windows.

- Fix `pg_dumpall` to work when a source database sets `default_transaction_read_only` via `ALTER DATABASE` (Kevin Grittner)

Previously, the generated script would fail during restore.

- Fix ecpg's processing of lists of variables declared `varchar` (Zoltán Böszörményi)
- Make `contrib/lo` defend against incorrect trigger definitions (Marc Cousin)
- Update timezone data files to `tzdata` release 2013h for DST law changes in Argentina, Brazil, Jordan, Libya, Liechtenstein, Morocco, and Palestine. Also, new timezone abbreviations WIB, WIT, WITA for Indonesia.

E.144. Release 9.0.14

Release date: 2013-10-10

This release contains a variety of fixes from 9.0.13. For information about new features in the 9.0 major release, see Section E.158.

E.144.1. Migration to Version 9.0.14

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.13, see Section E.145.

E.144.2. Changes

- Prevent corruption of multi-byte characters when attempting to case-fold identifiers (Andrew Dunstan)

  PostgreSQL case-folds non-ASCII characters only when using a single-byte server encoding.

- Fix checkpoint memory leak in background writer when `wal_level = hot_standby` (Naoya Anzai)

- Fix memory leak caused by `lo_open()` failure (Heikki Linnakangas)

- Fix memory overcommit bug when `work_mem` is using more than 24GB of memory (Stephen Frost)

- Fix deadlock bug in `libpq` when using SSL (Stephen Frost)

- Fix possible SSL state corruption in threaded `libpq` applications (Nick Phillips, Stephen Frost)

- Properly compute row estimates for boolean columns containing many NULL values (Andrew Gierth)
Previously tests like `col IS NOT TRUE` and `col IS NOT FALSE` did not properly factor in NULL values when estimating plan costs.

• Prevent pushing down WHERE clauses into unsafe UNION/INTERSECT subqueries (Tom Lane)

Subqueries of a UNION or INTERSECT that contain set-returning functions or volatile functions in their SELECT lists could be improperly optimized, leading to run-time errors or incorrect query results.

• Fix rare case of “failed to locate grouping columns” planner failure (Tom Lane)

• Improve view dumping code's handling of dropped columns in referenced tables (Tom Lane)

• Properly record index comments created using UNIQUE and PRIMARY KEY syntax (Andres Freund)

This fixes a parallel pg_restore failure.

• Fix REINDEX TABLE and REINDEX DATABASE to properly revalidate constraints and mark invalidated indexes as valid (Noah Misch)

REINDEX INDEX has always worked properly.

• Fix possible deadlock during concurrent CREATE INDEX CONCURRENTLY operations (Tom Lane)

• Fix regexp_matches() handling of zero-length matches (Jeevan Chalke)

Previously, zero-length matches like `'^'` could return too many matches.

• Fix crash for overly-complex regular expressions (Heikki Linnakangas)

• Fix regular expression match failures for back references combined with non-greedy quantifiers (Jeevan Chalke)

• Prevent CREATE FUNCTION from checking SET variables unless function body checking is enabled (Tom Lane)

• Allow ALTER DEFAULT PRIVILEGES to operate on schemas without requiring CREATE permission (Tom Lane)

• Loosen restriction on keywords used in queries (Tom Lane)

Specifically, lessen keyword restrictions for role names, language names, EXPLAIN and COPY options, and SET values. This allows COPY ... (FORMAT BINARY) to work as expected; previously BINARY needed to be quoted.

• Fix pgp_pub_decrypt() so it works for secret keys with passwords (Marko Kreen)

• Remove rare inaccurate warning during vacuum of index-less tables (Heikki Linnakangas)

• Ensure that VACUUM ANALYZE still runs the ANALYZE phase if its attempt to truncate the file is cancelled due to lock conflicts (Kevin Grittner)

• Avoid possible failure when performing transaction control commands (e.g ROLLBACK) in prepared queries (Tom Lane)

• Ensure that floating-point data input accepts standard spellings of “infinity” on all platforms (Tom Lane)

The C99 standard says that allowable spellings are inf, +inf, -inf, infinity, +infinity, and -infinity. Make sure we recognize these even if the platform's strtod function doesn't.
• Expand ability to compare rows to records and arrays (Rafal Rzepecki, Tom Lane)

• Update time zone data files to tzdata release 2013d for DST law changes in Israel, Morocco, Palestine, and Paraguay. Also, historical zone data corrections for Macquarie Island.

E.145. Release 9.0.13

Release date: 2013-04-04

This release contains a variety of fixes from 9.0.12. For information about new features in the 9.0 major release, see Section E.158.

E.145.1. Migration to Version 9.0.13

A dump/restore is not required for those running 9.0.X.

However, this release corrects several errors in management of GiST indexes. After installing this update, it is advisable to REINDEX any GiST indexes that meet one or more of the conditions described below.

Also, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.145.2. Changes

• Fix insecure parsing of server command-line switches (Mitsumasa Kondo, Kyotaro Horiguchi)

A connection request containing a database name that begins with “-” could be crafted to damage or destroy files within the server's data directory, even if the request is eventually rejected. (CVE-2013-1899)

• Reset OpenSSL randomness state in each postmaster child process (Marko Kreen)

This avoids a scenario wherein random numbers generated by contrib/pgcrypto functions might be relatively easy for another database user to guess. The risk is only significant when the postmaster is configured with ssl = on but most connections don't use SSL encryption. (CVE-2013-1900)

• Fix GiST indexes to not use “fuzzy” geometric comparisons when it's not appropriate to do so (Alexander Korotkov)

The core geometric types perform comparisons using “fuzzy” equality, but gist_box_same must do exact comparisons, else GiST indexes using it might become inconsistent. After installing this update, users should REINDEX any GiST indexes on box, polygon, circle, or point columns, since all of these use gist_box_same.

• Fix erroneous range-union and penalty logic in GiST indexes that use contrib/btree_gist for variable-width data types, that is text, bytea, bit, and numeric columns (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in useless index bloat. Users are advised to REINDEX such indexes after installing this update.

• Fix bugs in GiST page splitting code for multi-column indexes (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in indexes that are unnecessarily inefficient to search. Users are advised to REINDEX multi-column GiST indexes after installing this update.

• Fix gist_point_consistent to handle fuzziness consistently (Alexander Korotkov)
Index scans on GiST indexes on point columns would sometimes yield results different from a sequential scan, because gist_point_consistent disagreed with the underlying operator code about whether to do comparisons exactly or fuzzily.

• Fix buffer leak in WAL replay (Heikki Linnakangas)

This bug could result in “incorrect local pin count” errors during replay, making recovery impossible.

• Fix race condition in DELETE RETURNING (Tom Lane)

Under the right circumstances, DELETE RETURNING could attempt to fetch data from a shared buffer that the current process no longer has any pin on. If some other process changed the buffer meanwhile, this would lead to garbage RETURNING output, or even a crash.

• Fix infinite-loop risk in regular expression compilation (Tom Lane, Don Porter)

• Fix potential null-pointer dereference in regular expression compilation (Tom Lane)

• Fix to_char() to use ASCII-only case-folding rules where appropriate (Tom Lane)

This fixes misbehavior of some template patterns that should be locale-independent, but mishandled “I” and “i” in Turkish locales.

• Fix unwanted rejection of timestamp 1999-12-31 24:00:00 (Tom Lane)

• Fix logic error when a single transaction does UNLISTEN then LISTEN (Tom Lane)

The session wound up not listening for notify events at all, though it surely should listen in this case.

• Remove useless “picksplit doesn’t support secondary split” log messages (Josh Hansen, Tom Lane)

This message seems to have been added in expectation of code that was never written, and probably never will be, since GiST’s default handling of secondary splits is actually pretty good. So stop nagging end users about it.

• Fix possible failure to send a session’s last few transaction commit/abort counts to the statistics collector (Tom Lane)

• Eliminate memory leaks in PL/Perl's spi_prepare() function (Alex Hunsaker, Tom Lane)

• Fix pg_dumpall to handle database names containing “=” correctly (Heikki Linnakangas)

• Avoid crash in pg_dump when an incorrect connection string is given (Heikki Linnakangas)

• Ignore invalid indexes in pg_dump and pg_upgrade (Michael Paquier, Bruce Momijian)

Dumping invalid indexes can cause problems at restore time, for example if the reason the index creation failed was because it tried to enforce a uniqueness condition not satisfied by the table's data. Also, if the index creation is in fact still in progress, it seems reasonable to consider it to be an uncommitted DDL change, which pg_dump wouldn't be expected to dump anyway. pg_upgrade now also skips invalid indexes rather than failing.

• Fix contrib/pg_trgm's similarity() function to return zero for trigram-less strings (Tom Lane)

Previously it returned NaN due to internal division by zero.

• Update time zone data files to tzdata release 2013b for DST law changes in Chile, Haiti, Morocco, Paraguay, and some Russian areas. Also, historical zone data corrections for numerous places.
Also, update the time zone abbreviation files for recent changes in Russia and elsewhere: CHOT, GET, IRKT, KGT, KRAT, MAGT, MAWT, MSK, NOVT, OMST, TKT, VLAT, WST, YAKT, YEKT now follow their current meanings, and VOLT (Europe/Volgograd) and MIST (Antarctica/Macquarie) are added to the default abbreviations list.

E.146. Release 9.0.12

Release date: 2013-02-07

This release contains a variety of fixes from 9.0.11. For information about new features in the 9.0 major release, see Section E.158.

E.146.1. Migration to Version 9.0.12

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.146.2. Changes

• Prevent execution of `enum_recv` from SQL (Tom Lane)

  The function was misdeclared, allowing a simple SQL command to crash the server. In principle an attacker might be able to use it to examine the contents of server memory. Our thanks to Sumit Soni (via Secunia SVCRP) for reporting this issue. (CVE-2013-0255)

• Fix multiple problems in detection of when a consistent database state has been reached during WAL replay (Fujii Masao, Heikki Linnakangas, Simon Riggs, Andres Freund)

• Update minimum recovery point when truncating a relation file (Heikki Linnakangas)

  Once data has been discarded, it’s no longer safe to stop recovery at an earlier point in the timeline.

• Fix missing cancellations in hot standby mode (Noah Misch, Simon Riggs)

  The need to cancel conflicting hot-standby queries would sometimes be missed, allowing those queries to see inconsistent data.

• Fix SQL grammar to allow subscripting or field selection from a sub-SELECT result (Tom Lane)

• Fix performance problems with autovacuum truncation in busy workloads (Jan Wieck)

  Truncation of empty pages at the end of a table requires exclusive lock, but autovacuum was coded to fail (and release the table lock) when there are conflicting lock requests. Under load, it is easily possible that truncation would never occur, resulting in table bloat. Fix by performing a partial truncation, releasing the lock, then attempting to re-acquire the lock and continue. This fix also greatly reduces the average time before autovacuum releases the lock after a conflicting request arrives.

• Protect against race conditions when scanning `pg_tablespace` (Stephen Frost, Tom Lane)

  `CREATE DATABASE` and `DROP DATABASE` could misbehave if there were concurrent updates of `pg_tablespace` entries.

• Prevent `DROP OWNED` from trying to drop whole databases or tablespaces (Álvaro Herrera)

  For safety, ownership of these objects must be reassigned, not dropped.

• Fix error in `vacuum_freeze_table_age` implementation (Andres Freund)
In installations that have existed for more than vacuum_freeze_min_age transactions, this mistake prevented autovacuum from using partial-table scans, so that a full-table scan would always happen instead.

- Prevent misbehavior when a RowExpr or XmlExpr is parse-analyzed twice (Andres Freund, Tom Lane)

  This mistake could be user-visible in contexts such as CREATE TABLE LIKE INCLUDING INDEXES.

- Improve defenses against integer overflow in hashtable sizing calculations (Jeff Davis)

- Reject out-of-range dates in to_date() (Hitoshi Harada)

- Ensure that non-ASCII prompt strings are translated to the correct code page on Windows (Alexander Law, Noah Misch)

  This bug affected psql and some other client programs.

- Fix possible crash in psql's \? command when not connected to a database (Meng Qingzhong)

- Fix pg_upgrade to deal with invalid indexes safely (Bruce Momjian)

- Fix one-byte buffer overrun in libpq's PQprintTuples (Xi Wang)

  This ancient function is not used anywhere by PostgreSQL itself, but it might still be used by some client code.

- Make ecpglib use translated messages properly (Chen Huajun)

- Properly install ecpg_compat and pgtypes libraries on MSVC (Jiang Guiqing)

- Include our version of isinf() in libecpg if it's not provided by the system (Jiang Guiqing)

- Rearrange configure's tests for supplied functions so it is not fooled by bogus exports from libedit/libreadline (Christoph Berg)

- Ensure Windows build number increases over time (Magnus Hagander)

- Make pgxs build executables with the right .exe suffix when cross-compiling for Windows (Zoltan Boszormenyi)

- Add new timezone abbreviation FET (Tom Lane)

  This is now used in some eastern-European time zones.

### E.147. Release 9.0.11

**Release date:** 2012-12-06

This release contains a variety of fixes from 9.0.10. For information about new features in the 9.0 major release, see Section E.158.

### E.147.1. Migration to Version 9.0.11

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.
E.147.2. Changes

- Fix multiple bugs associated with CREATE INDEX CONCURRENTLY (Andres Freund, Tom Lane)

  Fix CREATE INDEX CONCURRENTLY to use in-place updates when changing the state of an index's pg_index row. This prevents race conditions that could cause concurrent sessions to miss updating the target index, thus resulting in corrupt concurrently-created indexes.

  Also, fix various other operations to ensure that they ignore invalid indexes resulting from a failed CREATE INDEX CONCURRENTLY command. The most important of these is VACUUM, because an auto-vacuum could easily be launched on the table before corrective action can be taken to fix or remove the invalid index.

- Fix buffer locking during WAL replay (Tom Lane)

  The WAL replay code was insufficiently careful about locking buffers when replaying WAL records that affect more than one page. This could result in hot standby queries transiently seeing inconsistent states, resulting in wrong answers or unexpected failures.

- Fix an error in WAL generation logic for GIN indexes (Tom Lane)

  This could result in index corruption, if a torn-page failure occurred.

- Properly remove startup process's virtual XID lock when promoting a hot standby server to normal running (Simon Riggs)

  This oversight could prevent subsequent execution of certain operations such as CREATE INDEX CONCURRENTLY.

- Avoid bogus “out-of-sequence timeline ID” errors in standby mode (Heikki Linnakangas)

- Prevent the postmaster from launching new child processes after it's received a shutdown signal (Tom Lane)

  This mistake could result in shutdown taking longer than it should, or even never completing at all without additional user action.

- Avoid corruption of internal hash tables when out of memory (Hitoshi Harada)

- Fix planning of non-strict equivalence clauses above outer joins (Tom Lane)

  The planner could derive incorrect constraints from a clause equating a non-strict construct to something else, for example WHERE COALESCE(foo, 0) = 0 when foo is coming from the nullable side of an outer join.

- Improve planner's ability to prove exclusion constraints from equivalence classes (Tom Lane)

- Fix partial-row matching in hashed subplans to handle cross-type cases correctly (Tom Lane)

  This affects multicolumn NOT IN subplans, such as WHERE (a, b) NOT IN (SELECT x, y FROM ...) when for instance b and y are int4 and int8 respectively. This mistake led to wrong answers or crashes depending on the specific datatypes involved.

- Acquire buffer lock when re-fetching the old tuple for an AFTER ROW UPDATE/DELETE trigger (Andres Freund)

  In very unusual circumstances, this oversight could result in passing incorrect data to the precheck logic for a foreign-key enforcement trigger. That could result in a crash, or in an incorrect decision about whether to fire the trigger.

- Fix ALTER COLUMN TYPE to handle inherited check constraints properly (Pavan Deolasee)
This worked correctly in pre-8.4 releases, and now works correctly in 8.4 and later.

- Fix REASSIGN OWNED to handle grants on tablespaces (Álvaro Herrera)
- Ignore incorrect `pg_attribute` entries for system columns for views (Tom Lane)
  Views do not have any system columns. However, we forgot to remove such entries when converting a table to a view. That's fixed properly for 9.3 and later, but in previous branches we need to defend against existing mis-converted views.
- Fix rule printing to dump `INSERT INTO table DEFAULT VALUES` correctly (Tom Lane)
- Guard against stack overflow when there are too many UNION/INTERSECT/EXCEPT clauses in a query (Tom Lane)
- Prevent platform-dependent failures when dividing the minimum possible integer value by -1 (Xi Wang, Tom Lane)
- Fix possible access past end of string in date parsing (Hitoshi Harada)
- Fix failure to advance XID epoch if XID wraparound happens during a checkpoint and `wal_level` is `hot_standby` (Tom Lane, Andres Freund)
  While this mistake had no particular impact on PostgreSQL itself, it was bad for applications that rely on `txid_current()` and related functions: the TXID value would appear to go backwards.
- Produce an understandable error message if the length of the path name for a Unix-domain socket exceeds the platform-specific limit (Tom Lane, Andrew Dunstan)
  Formerly, this would result in something quite unhelpful, such as “Non-recoverable failure in name resolution”.
- Fix memory leaks when sending composite column values to the client (Tom Lane)
- Make `pg_ctl` more robust about reading the `postmaster.pid` file (Heikki Linnakangas)
  Fix race conditions and possible file descriptor leakage.
- Fix possible crash in `psql` if incorrectly-encoded data is presented and the `client_encoding` setting is a client-only encoding, such as SJIS (Jiang Guiqing)
- Fix bugs in the `restore.sql` script emitted by `pg_dump` in `tar` output format (Tom Lane)
  The script would fail outright on tables whose names include upper-case characters. Also, make the script capable of restoring data in `--inserts` mode as well as the regular `COPY` mode.
- Fix `pg_restore` to accept POSIX-conformant `tar` files (Brian Weaver, Tom Lane)
  The original coding of `pg_dump`'s `tar` output mode produced files that are not fully conformant with the POSIX standard. This has been corrected for version 9.3. This patch updates previous branches so that they will accept both the incorrect and the corrected formats, in hopes of avoiding compatibility problems when 9.3 comes out.
- Fix `pg_resetxlog` to locate `postmaster.pid` correctly when given a relative path to the data directory (Tom Lane)
  This mistake could lead to `pg_resetxlog` not noticing that there is an active postmaster using the data directory.
- Fix `libpq`'s `lo_import()` and `lo_export()` functions to report file I/O errors properly (Tom Lane)
• Fix ecpg's processing of nested structure pointer variables (Muhammad Usama)
• Fix ecpg's ecpg_get_data function to handle arrays properly (Michael Meskes)
• Make contrib/pageinspect's btree page inspection functions take buffer locks while examining pages (Tom Lane)
• Fix pgxs support for building loadable modules on AIX (Tom Lane)
  Building modules outside the original source tree didn't work on AIX.
• Update time zone data files to tzdata release 2012j for DST law changes in Cuba, Israel, Jordan, Libya, Palestine, Western Samoa, and portions of Brazil.

E.148. Release 9.0.10

Release date: 2012-09-24

This release contains a variety of fixes from 9.0.9. For information about new features in the 9.0 major release, see Section E.158.

E.148.1. Migration to Version 9.0.10

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.148.2. Changes

• Fix planner's assignment of executor parameters, and fix executor's rescan logic for CTE plan nodes (Tom Lane)
  These errors could result in wrong answers from queries that scan the same WITH subquery multiple times.
• Improve page-splitting decisions in GiST indexes (Alexander Korotkov, Robert Haas, Tom Lane)
  Multi-column GiST indexes might suffer unexpected bloat due to this error.
• Fix cascading privilege revoke to stop if privileges are still held (Tom Lane)
  If we revoke a grant option from some role X, but X still holds that option via a grant from someone else, we should not recursively revoke the corresponding privilege from role(s) Y that X had granted it to.
• Improve error messages for Hot Standby misconfiguration errors (Gurjeet Singh)
• Fix handling of SIGFPE when PL/Perl is in use (Andres Freund)
  Perl resets the process's SIGFPE handler to SIG_IGN, which could result in crashes later on. Restore the normal Postgres signal handler after initializing PL/Perl.
• Prevent PL/Perl from crashing if a recursive PL/Perl function is redefined while being executed (Tom Lane)
• Work around possible misoptimization in PL/Perl (Tom Lane)
  Some Linux distributions contain an incorrect version of pthread.h that results in incorrect compiled code in PL/Perl, leading to crashes if a PL/Perl function calls another one that throws an error.
• Fix pg_upgrade's handling of line endings on Windows (Andrew Dunstan)

Previously, pg_upgrade might add or remove carriage returns in places such as function bodies.

• On Windows, make pg_upgrade use backslash path separators in the scripts it emits (Andrew Dunstan)

• Update time zone data files to tzdata release 2012f for DST law changes in Fiji

E.149. Release 9.0.9

Release date: 2012-08-17

This release contains a variety of fixes from 9.0.8. For information about new features in the 9.0 major release, see Section E.158.

E.149.1. Migration to Version 9.0.9

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.149.2. Changes

• Prevent access to external files/URLs via XML entity references (Noah Misch, Tom Lane)

xml_parse() would attempt to fetch external files or URLs as needed to resolve DTD and entity references in an XML value, thus allowing unprivileged database users to attempt to fetch data with the privileges of the database server. While the external data wouldn't get returned directly to the user, portions of it could be exposed in error messages if the data didn't parse as valid XML; and in any case the mere ability to check existence of a file might be useful to an attacker. (CVE-2012-3489)

• Prevent access to external files/URLs via contrib/xml2's xslt_process() (Peter Eisenbraut)

libxslt offers the ability to read and write both files and URLs through stylesheet commands, thus allowing unprivileged database users to both read and write data with the privileges of the database server. Disable that through proper use of libxslt's security options. (CVE-2012-3488)

Also, remove xslt_process()'s ability to fetch documents and stylesheets from external files/URLs. While this was a documented "feature", it was long regarded as a bad idea. The fix for CVE-2012-3489 broke that capability, and rather than expend effort on trying to fix it, we're just going to summarily remove it.

• Prevent too-early recycling of btree index pages (Noah Misch)

When we allowed read-only transactions to skip assigning XIDs, we introduced the possibility that a deleted btree page could be recycled while a read-only transaction was still in flight to it. This would result in incorrect index search results. The probability of such an error occurring in the field seems very low because of the timing requirements, but nonetheless it should be fixed.

• Fix crash-safety bug with newly-created-or-reset sequences (Tom Lane)

If ALTER SEQUENCE was executed on a freshly created or reset sequence, and then precisely one nextval() call was made on it, and then the server crashed, WAL replay would restore the sequence to a state in which it appeared that no nextval() had been done, thus allowing the first sequence value to be returned again by the next nextval() call. In particular this could manifest
for serial columns, since creation of a serial column's sequence includes an ALTER SEQUENCE
OWNED BY step.

- Fix txid_current() to report the correct epoch when not in hot standby (Heikki Linnakangas)

  This fixes a regression introduced in the previous minor release.

- Fix bug in startup of Hot Standby when a master transaction has many subtransactions (Andres
  Freund)

  This mistake led to failures reported as “out-of-order XID insertion in KnownAssignedXids”.

- Ensure the backup_label file is fsync'd after pg_start_backup() (Dave Kerr)

- Fix timeout handling in walsender processes (Tom Lane)

  WAL sender background processes neglected to establish a SIGALRM handler, meaning they
  would wait forever in some corner cases where a timeout ought to happen.

- Back-patch 9.1 improvement to compress the fsync request queue (Robert Haas)

  This improves performance during checkpoints. The 9.1 change has now seen enough field testing
  to seem safe to back-patch.

- Fix LISTEN/NOTIFY to cope better with I/O problems, such as out of disk space (Tom Lane)

  After a write failure, all subsequent attempts to send more NOTIFY messages would fail with mes-
  sages like “Could not read from file “pg_notify/nnnn” at offset nnnnn: Success”.

- Only allow autovacuum to be auto-canceled by a directly blocked process (Tom Lane)

  The original coding could allow inconsistent behavior in some cases; in particular, an autovacuum
  could get canceled after less than deadlock_timeout grace period.

- Improve logging of autovacuum cancels (Robert Haas)

- Fix log collector so that log_truncate_on_rotation works during the very first log rotation
  after server start (Tom Lane)

- Fix WITH attached to a nested set operation (UNION/INTERSECT/EXCEPT) (Tom Lane)

- Ensure that a whole-row reference to a subquery doesn't include any extra GROUP BY or ORDER
  BY columns (Tom Lane)

- Disallow copying whole-row references in CHECK constraints and index definitions during CREATE
  TABLE (Tom Lane)

  This situation can arise in CREATE TABLE with LIKE or INHERITS. The copied whole-row
  variable was incorrectly labeled with the row type of the original table not the new one. Rejecting
  the case seems reasonable for LIKE, since the row types might well diverge later. For INHERITS
  we should ideally allow it, with an implicit coercion to the parent table's row type; but that will
  require more work than seems safe to back-patch.

- Fix memory leak in ARRAY(SELECT ...) subqueries (Heikki Linnakangas, Tom Lane)

- Fix extraction of common prefixes from regular expressions (Tom Lane)

  The code could get confused by quantified parenthesized subexpressions, such as ^{foo}?bar.
  This would lead to incorrect index optimization of searches for such patterns.

- Fix bugs with parsing signed hh:mm and hh:mm:ss fields in interval constants (Amit Kapila,
  Tom Lane)
Release Notes

• Use Postgres’ encoding conversion functions, not Python’s, when converting a Python Unicode string to the server encoding in PL/Python (Jan Urbanski)

  This avoids some corner-case problems, notably that Python doesn’t support all the encodings Postgres does. A notable functional change is that if the server encoding is SQL_ASCII, you will get the UTF-8 representation of the string; formerly, any non-ASCII characters in the string would result in an error.

• Fix mapping of PostgreSQL encodings to Python encodings in PL/Python (Jan Urbanski)

• Report errors properly in contrib/xml2\xslt_process() (Tom Lane)

• Update time zone data files to tzdata release 2012e for DST law changes in Morocco and Tokelau

E.150. Release 9.0.8

Release date: 2012-06-04

This release contains a variety of fixes from 9.0.7. For information about new features in the 9.0 major release, see Section E.158.

E.150.1. Migration to Version 9.0.8

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.150.2. Changes

• Fix incorrect password transformation in contrib/pgcrypto's DES crypt() function (Solar Designer)

  If a password string contained the byte value 0x80, the remainder of the password was ignored, causing the password to be much weaker than it appeared. With this fix, the rest of the string is properly included in the DES hash. Any stored password values that are affected by this bug will thus no longer match, so the stored values may need to be updated. (CVE-2012-2143)

• Ignore SECURITY DEFINER and SET attributes for a procedural language's call handler (Tom Lane)

  Applying such attributes to a call handler could crash the server. (CVE-2012-2655)

• Allow numeric timezone offsets in timestamp input to be up to 16 hours away from UTC (Tom Lane)

  Some historical time zones have offsets larger than 15 hours, the previous limit. This could result in dumped data values being rejected during reload.

• Fix timestamp conversion to cope when the given time is exactly the last DST transition time for the current timezone (Tom Lane)

  This oversight has been there a long time, but was not noticed previously because most DST-using zones are presumed to have an indefinite sequence of future DST transitions.

• Fix text to name and char to name casts to perform string truncation correctly in multibyte encodings (Karl Schnaitter)

• Fix memory copying bug in to_tsquery() (Heikki Linnakangas)
• Ensure `txid_current()` reports the correct epoch when executed in hot standby (Simon Riggs)

• Fix planner's handling of outer PlaceHolderVars within subqueries (Tom Lane)

This bug concerns sub-SELECTs that reference variables coming from the nullable side of an outer join of the surrounding query. In 9.1, queries affected by this bug would fail with “ERROR: Upper-level PlaceHolderVar found where not expected”. But in 9.0 and 8.4, you'd silently get possibly-wrong answers, since the value transmitted into the subquery wouldn't go to null when it should.

• Fix slow session startup when `pg_attribute` is very large (Tom Lane)

If `pg_attribute` exceeds one-fourth of `shared_buffers`, cache rebuilding code that is sometimes needed during session start would trigger the synchronized-scan logic, causing it to take many times longer than normal. The problem was particularly acute if many new sessions were starting at once.

• Ensure sequential scans check for query cancel reasonably often (Merlin Moncure)

A scan encountering many consecutive pages that contain no live tuples would not respond to interrupts meanwhile.

• Ensure the Windows implementation of `PGSemaphoreLock()` clears `ImmediateInterruptOK` before returning (Tom Lane)

This oversight meant that a query-cancel interrupt received later in the same query could be accepted at an unsafe time, with unpredictable but not good consequences.

• Show whole-row variables safely when printing views or rules (Abbas Butt, Tom Lane)

Corner cases involving ambiguous names (that is, the name could be either a table or column name of the query) were printed in an ambiguous way, risking that the view or rule would be interpreted differently after dump and reload. Avoid the ambiguous case by attaching a no-op cast.

• Fix `COPY FROM` to properly handle null marker strings that correspond to invalid encoding (Tom Lane)

A null marker string such as `E'\0'` should work, and did work in the past, but the case got broken in 8.4.

• Ensure autovacuum worker processes perform stack depth checking properly (Heikki Linnakangas)

Previously, infinite recursion in a function invoked by auto-ANALYZE could crash worker processes.

• Fix logging collector to not lose log coherency under high load (Andrew Dunstan)

The collector previously could fail to reassemble large messages if it got too busy.

• Fix logging collector to ensure it will restart file rotation after receiving SIGHUP (Tom Lane)

• Fix WAL replay logic for GIN indexes to not fail if the index was subsequently dropped (Tom Lane)

• Fix memory leak in PL/pgSQL's `RETURN NEXT` command (Joe Conway)

• Fix PL/pgSQL's `GET DIAGNOSTICS` command when the target is the function's first variable (Tom Lane)

• Fix potential access off the end of memory in `psql`'s expanded display (`\x`) mode (Peter Eisentraut)

• Fix several performance problems in `pg_dump` when the database contains many objects (Jeff Janes, Tom Lane)
pg_dump could get very slow if the database contained many schemas, or if many objects are in dependency loops, or if there are many owned sequences.

- Fix pg_upgrade for the case that a database stored in a non-default tablespace contains a table in the cluster's default tablespace (Bruce Momjian)
- In ecpg, fix rare memory leaks and possible overwrite of one byte after the sqlca_t structure (Peter Eisentraut)
- Fix contrib/dblink's dblink_exec() to not leak temporary database connections upon error (Tom Lane)
- Fix contrib/dblink to report the correct connection name in error messages (Kyotaro Horiguchi)
- Fix contrib/vacuumlo to use multiple transactions when dropping many large objects (Tim Lewis, Robert Haas, Tom Lane)
  This change avoids exceeding max_locks_per_transaction when many objects need to be dropped. The behavior can be adjusted with the new -l (limit) option.
- Update time zone data files to tzdata release 2012c for DST law changes in Antarctica, Armenia, Chile, Cuba, Falkland Islands, Gaza, Haiti, Hebron, Morocco, Syria, and Tokelau Islands; also historical corrections for Canada.

E.151. Release 9.0.7

Release date: 2012-02-27

This release contains a variety of fixes from 9.0.6. For information about new features in the 9.0 major release, see Section E.158.

E.151.1. Migration to Version 9.0.7

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.6, see Section E.152.

E.151.2. Changes

- Require execute permission on the trigger function for CREATE TRIGGER (Robert Haas)
  This missing check could allow another user to execute a trigger function with forged input data, by installing it on a table he owns. This is only of significance for trigger functions marked SECURITY DEFINER, since otherwise trigger functions run as the table owner anyway. (CVE-2012-0866)
- Remove arbitrary limitation on length of common name in SSL certificates (Heikki Linnakangas)
  Both libpq and the server truncated the common name extracted from an SSL certificate at 32 bytes. Normally this would cause nothing worse than an unexpected verification failure, but there are some rather-implausible scenarios in which it might allow one certificate holder to impersonate another. The victim would have to have a common name exactly 32 bytes long, and the attacker would have to persuade a trusted CA to issue a certificate in which the common name has that string as a prefix. Impersonating a server would also require some additional exploit to redirect client connections. (CVE-2012-0867)
- Convert newlines to spaces in names written in pg_dump comments (Robert Haas)
pg_dump was incautious about sanitizing object names that are emitted within SQL comments in its output script. A name containing a newline would at least render the script syntactically incorrect. Maliciously crafted object names could present a SQL injection risk when the script is reloaded. (CVE-2012-0868)

• Fix btree index corruption from insertions concurrent with vacuuming (Tom Lane)

An index page split caused by an insertion could sometimes cause a concurrently-running VACUUM to miss removing index entries that it should remove. After the corresponding table rows are removed, the dangling index entries would cause errors (such as “could not read block N in file ...”) or worse, silently wrong query results after unrelated rows are re-inserted at the now-free table locations. This bug has been present since release 8.2, but occurs so infrequently that it was not diagnosed until now. If you have reason to suspect that it has happened in your database, reindexing the affected index will fix things.

• Fix transient zeroing of shared buffers during WAL replay (Tom Lane)

The replay logic would sometimes zero and refill a shared buffer, so that the contents were transiently invalid. In hot standby mode this can result in a query that's executing in parallel seeing garbage data. Various symptoms could result from that, but the most common one seems to be “invalid memory alloc request size”.

• Fix postmaster to attempt restart after a hot-standby crash (Tom Lane)

A logic error caused the postmaster to terminate, rather than attempt to restart the cluster, if any backend process crashed while operating in hot standby mode.

• Fix CLUSTER/VACUUM FULL handling of toast values owned by recently-updated rows (Tom Lane)

This oversight could lead to “duplicate key value violates unique constraint” errors being reported against the toast table's index during one of these commands.

• Update per-column permissions, not only per-table permissions, when changing table owner (Tom Lane)

Failure to do this meant that any previously granted column permissions were still shown as having been granted by the old owner. This meant that neither the new owner nor a superuser could revoke the now-untraceable-to-table-owner permissions.

• Support foreign data wrappers and foreign servers in REASSIGN OWNED (Alvaro Herrera)

This command failed with “unexpected classid” errors if it needed to change the ownership of any such objects.

• Allow non-existent values for some settings in ALTER USER/DATABASE SET (Heikki Linnakangas)

Allow default_text_search_config, default_tablespace, and temp_tablespaces to be set to names that are not known. This is because they might be known in another database where the setting is intended to be used, or for the tablespace cases because the tablespace might not be created yet. The same issue was previously recognized for search_path, and these settings now act like that one.

• Avoid crashing when we have problems deleting table files post-commit (Tom Lane)

Dropping a table should lead to deleting the underlying disk files only after the transaction commits. In event of failure then (for instance, because of wrong file permissions) the code is supposed to just emit a warning message and go on, since it's too late to abort the transaction. This logic got broken as of release 8.4, causing such situations to result in a PANIC and an unrestartable database.
• Recover from errors occurring during WAL replay of DROP TABLESPACE (Tom Lane)

Replay will attempt to remove the tablespace's directories, but there are various reasons why this might fail (for example, incorrect ownership or permissions on those directories). Formerly the replay code would panic, rendering the database unrestartable without manual intervention. It seems better to log the problem and continue, since the only consequence of failure to remove the directories is some wasted disk space.

• Fix race condition in logging AccessExclusiveLocks for hot standby (Simon Riggs)

Sometimes a lock would be logged as being held by “transaction zero”. This is at least known to produce assertion failures on slave servers, and might be the cause of more serious problems.

• Track the OID counter correctly during WAL replay, even when it wraps around (Tom Lane)

Previously the OID counter would remain stuck at a high value until the system exited replay mode. The practical consequences of that are usually nil, but there are scenarios wherein a standby server that's been promoted to master might take a long time to advance the OID counter to a reasonable value once values are needed.

• Prevent emitting misleading “consistent recovery state reached” log message at the beginning of crash recovery (Heikki Linnakangas)

• Fix initial value of pg_stat_replication.replay_location (Fujii Masao)

Previously, the value shown would be wrong until at least one WAL record had been replayed.

• Fix regular expression back-references with * attached (Tom Lane)

Rather than enforcing an exact string match, the code would effectively accept any string that satisfies the pattern sub-expression referenced by the back-reference symbol.

A similar problem still afflicts back-references that are embedded in a larger quantified expression, rather than being the immediate subject of the quantifier. This will be addressed in a future PostgreSQL release.

• Fix recently-introduced memory leak in processing of inet/cidr values (Heikki Linnakangas)

A patch in the December 2011 releases of PostgreSQL caused memory leakage in these operations, which could be significant in scenarios such as building a btree index on such a column.

• Fix dangling pointer after CREATE TABLE AS/SELECT INTO in a SQL-language function (Tom Lane)

In most cases this only led to an assertion failure in assert-enabled builds, but worse consequences seem possible.

• Avoid double close of file handle in syslogger on Windows (MauMau)

Ordinarily this error was invisible, but it would cause an exception when running on a debug version of Windows.

• Fix I/O-conversion-related memory leaks in plpgsql (Andres Freund, Jan Urbanski, Tom Lane)

Certain operations would leak memory until the end of the current function.

• Improve pg_dump's handling of inherited table columns (Tom Lane)

pg_dump mishandled situations where a child column has a different default expression than its parent column. If the default is textually identical to the parent's default, but not actually the same (for instance, because of schema search path differences) it would not be recognized as different, so
that after dump and restore the child would be allowed to inherit the parent's default. Child columns
that are NOT NULL where their parent is not could also be restored subtly incorrectly.

• Fix pg_restore's direct-to-database mode for INSERT-style table data (Tom Lane)

Direct-to-database restores from archive files made with --inserts or --column-inserts options fail when using pg_restore from a release dated September or December 2011, as a result of an oversight in a fix for another problem. The archive file itself is not at fault, and text-mode output is okay.

• Allow pg_upgrade to process tables containing regclass columns (Bruce Momjian)

Since pg_upgrade now takes care to preserve pg_class OIDs, there was no longer any reason for this restriction.

• Make libpq ignore ENOTDIR errors when looking for an SSL client certificate file (Magnus Hagander)

This allows SSL connections to be established, though without a certificate, even when the user's home directory is set to something like /dev/null.

• Fix some more field alignment issues in ecpg's SQLDA area (Zoltan Boszormenyi)

• Allow AT option in ecpg DEALLOCATE statements (Michael Meskes)

The infrastructure to support this has been there for awhile, but through an oversight there was still an error check rejecting the case.

• Do not use the variable name when defining a varchar structure in ecpg (Michael Meskes)

• Fix contrib/auto_explain's JSON output mode to produce valid JSON (Andrew Dunstan)

The output used brackets at the top level, when it should have used braces.

• Fix error in contrib/intarray's int[] & int[] operator (Guillaume Lelarge)

If the smallest integer the two input arrays have in common is 1, and there are smaller values in either array, then 1 would be incorrectly omitted from the result.

• Fix error detection in contrib/pgcrypto's encrypt_iv() and decrypt_iv() (Marko Kreen)

These functions failed to report certain types of invalid-input errors, and would instead return random garbage values for incorrect input.

• Fix one-byte buffer overrun in contrib/test_parser (Paul Guyot)

The code would try to read one more byte than it should, which would crash in corner cases. Since contrib/test_parser is only example code, this is not a security issue in itself, but bad example code is still bad.

• Use __sync_lock_test_and_set() for spinlocks on ARM, if available (Martin Pitt)

This function replaces our previous use of the SWPB instruction, which is deprecated and not available on ARMv6 and later. Reports suggest that the old code doesn't fail in an obvious way on recent ARM boards, but simply doesn't interlock concurrent accesses, leading to bizarre failures in multiprocess operation.

• Use -fexcess-precision=standard option when building with gcc versions that accept it (Andrew Dunstan)

This prevents assorted scenarios wherein recent versions of gcc will produce creative results.
• Allow use of threaded Python on FreeBSD (Chris Rees)

Our configure script previously believed that this combination wouldn't work; but FreeBSD fixed
the problem, so remove that error check.

E.152. Release 9.0.6

Release date: 2011-12-05

This release contains a variety of fixes from 9.0.5. For information about new features in the 9.0 major
release, see Section E.158.

E.152.1. Migration to Version 9.0.6

A dump/restore is not required for those running 9.0.X.

However, a longstanding error was discovered in the definition of the information_schema.referential_constraints view. If you rely on correct results from that
view, you should replace its definition as explained in the first changelog item below.

Also, if you are upgrading from a version earlier than 9.0.4, see Section E.154.

E.152.2. Changes

• Fix bugs in information_schema.referential_constraints view (Tom Lane)

This view was being insufficiently careful about matching the foreign-key constraint to the depend-
ed-on primary or unique key constraint. That could result in failure to show a foreign key constraint
at all, or showing it multiple times, or claiming that it depends on a different constraint than the
one it really does.

Since the view definition is installed by initdb, merely upgrading will not fix the prob-
lem. If you need to fix this in an existing installation, you can (as a superuser) drop
the information_schema schema then re-create it by sourcing SHAREDIR/informa-
tion_schema.sql. (Run pg_config --sharedir if you're uncertain where SHAREDIR
is.) This must be repeated in each database to be fixed.

• Fix possible crash during UPDATE or DELETE that joins to the output of a scalar-returning function
(Tom Lane)

A crash could only occur if the target row had been concurrently updated, so this problem surfaced
only intermittently.

• Fix incorrect replay of WAL records for GIN index updates (Tom Lane)

This could result in transiently failing to find index entries after a crash, or on a hot-standby server.
The problem would be repaired by the next VACUUM of the index, however.

• Fix TOAST-related data corruption during CREATE TABLE dest AS SELECT * FROM src
or INSERT INTO dest SELECT * FROM src (Tom Lane)

If a table has been modified by ALTER TABLE ADD COLUMN, attempts to copy its data verbatim
to another table could produce corrupt results in certain corner cases. The problem can only manifest
in this precise form in 8.4 and later, but we patched earlier versions as well in case there are other
code paths that could trigger the same bug.

• Fix possible failures during hot standby startup (Simon Riggs)

• Start hot standby faster when initial snapshot is incomplete (Simon Riggs)
Release Notes

- Fix race condition during toast table access from stale syscache entries (Tom Lane)

  The typical symptom was transient errors like “missing chunk number 0 for toast value NNNNN in pg_toast_2619”, where the cited toast table would always belong to a system catalog.

- Track dependencies of functions on items used in parameter default expressions (Tom Lane)

  Previously, a referenced object could be dropped without having dropped or modified the function, leading to misbehavior when the function was used. Note that merely installing this update will not fix the missing dependency entries; to do that, you’d need to CREATE OR REPLACE each such function afterwards. If you have functions whose defaults depend on non-built-in objects, doing so is recommended.

- Allow inlining of set-returning SQL functions with multiple OUT parameters (Tom Lane)

- Don’t trust deferred-unique indexes for join removal (Tom Lane and Marti Raudsepp)

  A deferred uniqueness constraint might not hold intra-transaction, so assuming that it does could give incorrect query results.

- Make DatumGetInetP() unpack inet datums that have a 1-byte header, and add a new macro, DatumGetInetPP(), that does not (Heikki Linnakangas)

  This change affects no core code, but might prevent crashes in add-on code that expects DatumGetInetP() to produce an unpacked datum as per usual convention.

- Improve locale support in money type’s input and output (Tom Lane)

  Aside from not supporting all standard lc_monetary formatting options, the input and output functions were inconsistent, meaning there were locales in which dumped money values could not be re-read.

- Don’t let transform_null_equals affect CASE foo WHEN NULL ... constructs (Heikki Linnakangas)

  transform_null_equals is only supposed to affect foo = NULL expressions written directly by the user, not equality checks generated internally by this form of CASE.

- Change foreign-key trigger creation order to better support self-referential foreign keys (Tom Lane)

  For a cascading foreign key that references its own table, a row update will fire both the ON UPDATE trigger and the CHECK trigger as one event. The ON UPDATE trigger must execute first, else the CHECK will check a non-final state of the row and possibly throw an inappropriate error. However, the firing order of these triggers is determined by their names, which generally sort in creation order since the triggers have auto-generated names following the convention “RI_ConstraintTrigger_NNNN”. A proper fix would require modifying that convention, which we will do in 9.2, but it seems risky to change it in existing releases. So this patch just changes the creation order of the triggers. Users encountering this type of error should drop and re-create the foreign key constraint to get its triggers into the right order.

- Avoid floating-point underflow while tracking buffer allocation rate (Greg Matthews)

  While harmless in itself, on certain platforms this would result in annoying kernel log messages.

- Preserve configuration file name and line number values when starting child processes under Windows (Tom Lane)

  Formerly, these would not be displayed correctly in the pg_settings view.

- Fix incorrect field alignment in ecpg’s SQLDA area (Zoltan Boszormenyi)

- Preserve blank lines within commands in psql’s command history (Robert Haas)
The former behavior could cause problems if an empty line was removed from within a string literal, for example.

- Fix pg_dump to dump user-defined casts between auto-generated types, such as table rowtypes (Tom Lane)

- Assorted fixes for pg_upgrade (Bruce Momjian)
  Handle exclusion constraints correctly, avoid failures on Windows, don't complain about mismatched toast table names in 8.4 databases.

- Use the preferred version of xsubpp to build PL/Perl, not necessarily the operating system's main copy (David Wheeler and Alex Hunsaker)

- Fix incorrect coding in contrib/dict_int and contrib/dict_xsyn (Tom Lane)
  Some functions incorrectly assumed that memory returned by palloc() is guaranteed zeroed.

- Fix assorted errors in contrib/unaccent's configuration file parsing (Tom Lane)

- Honor query cancel interrupts promptly in pgstatindex() (Robert Haas)

- Fix incorrect quoting of log file name in macOS start script (Sidar Lopez)

- Ensure VPATH builds properly install all server header files (Peter Eisentraut)

- Shorten file names reported in verbose error messages (Peter Eisentraut)
  Regular builds have always reported just the name of the C file containing the error message call, but VPATH builds formerly reported an absolute path name.

- Fix interpretation of Windows timezone names for Central America (Tom Lane)
  Map “Central America Standard Time” to CST6, not CST6CDT, because DST is generally not observed anywhere in Central America.

- Update time zone data files to tzdata release 2011n for DST law changes in Brazil, Cuba, Fiji, Palestine, Russia, and Samoa; also historical corrections for Alaska and British East Africa.

**E.153. Release 9.0.5**

**Release date:** 2011-09-26

This release contains a variety of fixes from 9.0.4. For information about new features in the 9.0 major release, see Section E.158.

**E.153.1. Migration to Version 9.0.5**

A dump/restore is not required for those running 9.0.X.

However, if you are upgrading from a version earlier than 9.0.4, see Section E.154.

**E.153.2. Changes**

- Fix catalog cache invalidation after a VACUUM FULL or CLUSTER on a system catalog (Tom Lane)
  In some cases the relocation of a system catalog row to another place would not be recognized by concurrent server processes, allowing catalog corruption to occur if they then tried to update that row. The worst-case outcome could be as bad as complete loss of a table.
• Fix incorrect order of operations during sinval reset processing, and ensure that TOAST OIDs are preserved in system catalogs (Tom Lane)

These mistakes could lead to transient failures after a VACUUM FULL or CLUSTER on a system catalog.

• Fix bugs in indexing of in-doubt HOT-updated tuples (Tom Lane)

These bugs could result in index corruption after reindexing a system catalog. They are not believed to affect user indexes.

• Fix multiple bugs in GiST index page split processing (Heikki Linnakangas)

The probability of occurrence was low, but these could lead to index corruption.

• Fix possible buffer overrun in tsvector_concat() (Tom Lane)

The function could underestimate the amount of memory needed for its result, leading to server crashes.

• Fix crash in xml_recv when processing a “standalone” parameter (Tom Lane)

• Make pg_options_to_table return NULL for an option with no value (Tom Lane)

Previously such cases would result in a server crash.

• Avoid possibly accessing off the end of memory in ANALYZE and in SJIS-2004 encoding conversion (Noah Misch)

This fixes some very-low-probability server crash scenarios.

• Protect pg_stat_reset_shared() against NULL input (Magnus Hagander)

• Fix possible failure when a recovery conflict deadlock is detected within a sub-transaction (Tom Lane)

• Avoid spurious conflicts while recycling btree index pages during hot standby (Noah Misch, Simon Riggs)

• Shut down WAL receiver if it's still running at end of recovery (Heikki Linnakangas)

The postmaster formerly panicked in this situation, but it's actually a legitimate case.

• Fix race condition in relcache init file invalidation (Tom Lane)

There was a window wherein a new backend process could read a stale init file but miss the inval messages that would tell it the data is stale. The result would be bizarre failures in catalog accesses, typically “could not read block 0 in file ...” later during startup.

• Fix memory leak at end of a GiST index scan (Tom Lane)

Commands that perform many separate GiST index scans, such as verification of a new GiST-based exclusion constraint on a table already containing many rows, could transiently require large amounts of memory due to this leak.

• Fix memory leak when encoding conversion has to be done on incoming command strings and LISTEN is active (Tom Lane)

• Fix incorrect memory accounting (leading to possible memory bloat) in tuplestores supporting holdable cursors and plpgsql's RETURN NEXT command (Tom Lane)

• Fix trigger WHEN conditions when both BEFORE and AFTER triggers exist (Tom Lane)
Evaluation of WHEN conditions for AFTER ROW UPDATE triggers could crash if there had been a BEFORE ROW trigger fired for the same update.

- Fix performance problem when constructing a large, lossy bitmap (Tom Lane)
- Fix join selectivity estimation for unique columns (Tom Lane)
  This fixes an erroneous planner heuristic that could lead to poor estimates of the result size of a join.
- Fix nested PlaceHolderVar expressions that appear only in sub-select target lists (Tom Lane)
  This mistake could result in outputs of an outer join incorrectly appearing as NULL.
- Allow the planner to assume that empty parent tables really are empty (Tom Lane)
  Normally an empty table is assumed to have a certain minimum size for planning purposes; but this heuristic seems to do more harm than good for the parent table of an inheritance hierarchy, which often is permanently empty.
- Allow nested EXISTS queries to be optimized properly (Tom Lane)
- Fix array- and path-creating functions to ensure padding bytes are zeroes (Tom Lane)
  This avoids some situations where the planner will think that semantically-equal constants are not equal, resulting in poor optimization.
- Fix EXPLAIN to handle gating Result nodes within inner-indexscan subplans (Tom Lane)
  The usual symptom of this oversight was “bogus varno” errors.
- Fix btree preprocessing of indexedcol IS NULL conditions (Dean Rasheed)
  Such a condition is unsatisfiable if combined with any other type of btree-indexable condition on the same index column. The case was handled incorrectly in 9.0.0 and later, leading to query output where there should be none.
- Work around gcc 4.6.0 bug that breaks WAL replay (Tom Lane)
  This could lead to loss of committed transactions after a server crash.
- Fix dump bug for VALUES in a view (Tom Lane)
- Disallow SELECT FOR UPDATE/SHARE on sequences (Tom Lane)
  This operation doesn't work as expected and can lead to failures.
- Fix VACUUM so that it always updates pg_class.reltuples/relpages (Tom Lane)
  This fixes some scenarios where autovacuum could make increasingly poor decisions about when to vacuum tables.
- Defend against integer overflow when computing size of a hash table (Tom Lane)
- Fix cases where CLUSTER might attempt to access already-removed TOAST data (Tom Lane)
- Fix premature timeout failures during initial authentication transaction (Tom Lane)
- Fix portability bugs in use of credentials control messages for “peer” authentication (Tom Lane)
- Fix SSPI login when multiple roundtrips are required (Ahmed Shinwari, Magnus Hagander)
  The typical symptom of this problem was “The function requested is not supported” errors during SSPI login.
• Fix failure when adding a new variable of a custom variable class to `postgresql.conf` (Tom Lane)

• Throw an error if `pg_hba.conf` contains `hostssl` but SSL is disabled (Tom Lane)
  This was concluded to be more user-friendly than the previous behavior of silently ignoring such lines.

• Fix failure when `DROP OWNED BY` attempts to remove default privileges on sequences (Shigeru Hanada)

• Fix typo in `pg_srand48` seed initialization (Andres Freund)
  This led to failure to use all bits of the provided seed. This function is not used on most platforms (only those without `srandom`), and the potential security exposure from a less-random-than-expected seed seems minimal in any case.

• Avoid integer overflow when the sum of `LIMIT` and `OFFSET` values exceeds $2^{63}$ (Heikki Linnakangas)

• Add overflow checks to `int4` and `int8` versions of `generate_series()` (Robert Haas)

• Fix trailing-zero removal in `to_char()` (Marti Raudsepp)
  In a format with `FM` and no digit positions after the decimal point, zeroes to the left of the decimal point could be removed incorrectly.

• Fix `pg_size_pretty()` to avoid overflow for inputs close to $2^{63}$ (Tom Lane)

• Weaken `plpgsql`'s check for typmod matching in record values (Tom Lane)
  An overly enthusiastic check could lead to discarding length modifiers that should have been kept.

• Correctly handle quotes in locale names during `initdb` (Heikki Linnakangas)
  The case can arise with some Windows locales, such as “People’s Republic of China”.

• In `pg_upgrade`, avoid dumping orphaned temporary tables (Bruce Momjian)
  This prevents situations wherein table OID assignments could get out of sync between old and new installations.

• Fix `pg_upgrade` to preserve toast tables' `refrozenxids` during an upgrade from 8.3 (Bruce Momjian)
  Failure to do this could lead to `pg_clog` files being removed too soon after the upgrade.

• In `pg_upgrade`, fix the `-l` (log) option to work on Windows (Bruce Momjian)

• In `pg_ctl`, support silent mode for service registrations on Windows (MauMau)

• Fix `psql`'s counting of script file line numbers during `COPY` from a different file (Tom Lane)

• Fix `pg_restore`'s direct-to-database mode for `standard_conforming_strings` (Tom Lane)
  `pg_restore` could emit incorrect commands when restoring directly to a database server from an archive file that had been made with `standard_conforming_strings` set to on.

• Be more user-friendly about unsupported cases for parallel `pg_restore` (Tom Lane)
  This change ensures that such cases are detected and reported before any restore actions have been taken.

• Fix write-past-buffer-end and memory leak in `libpq`'s LDAP service lookup code (Albe Laurenz)
• In libpq, avoid failures when using nonblocking I/O and an SSL connection (Martin Pihlak, Tom Lane)

• Improve libpq's handling of failures during connection startup (Tom Lane)

  In particular, the response to a server report of fork() failure during SSL connection startup is now saner.

• Improve libpq's error reporting for SSL failures (Tom Lane)

• Fix PQsetvalue() to avoid possible crash when adding a new tuple to a PGresult originally obtained from a server query (Andrew Chernow)

• Make ecpglib write double values with 15 digits precision (Akira Kurosawa)

• In ecpglib, be sure LC_NUMERIC setting is restored after an error (Michael Meskes)

• Apply upstream fix for blowfish signed-character bug (CVE-2011-2483) (Tom Lane)

  contrib/pg_crypto's blowfish encryption code could give wrong results on platforms where char is signed (which is most), leading to encrypted passwords being weaker than they should be.

• Fix memory leak in contrib/seg (Heikki Linnakangas)

• Fix pgstatindex() to give consistent results for empty indexes (Tom Lane)

• Allow building with perl 5.14 (Alex Hunsaker)

• Fix assorted issues with build and install file paths containing spaces (Tom Lane)

• Update time zone data files to tzdata release 2011i for DST law changes in Canada, Egypt, Russia, Samoa, and South Sudan.

E.154. Release 9.0.4

Release date: 2011-04-18

This release contains a variety of fixes from 9.0.3. For information about new features in the 9.0 major release, see Section E.158.

E.154.1. Migration to Version 9.0.4

A dump/restore is not required for those running 9.0.X.

However, if your installation was upgraded from a previous major release by running pg_upgrade, you should take action to prevent possible data loss due to a now-fixed bug in pg_upgrade. The recommended solution is to run VACUUM FREEZE on all TOAST tables. More information is available at http://wiki.postgresql.org/wiki/20110408pg_upgrade_fix.

E.154.2. Changes

• Fix pg_upgrade's handling of TOAST tables (Bruce Momjian)

  The pg_class.relrofxid value for TOAST tables was not correctly copied into the new installation during pg_upgrade. This could later result in pg_clog files being discarded while they were still needed to validate tuples in the TOAST tables, leading to “could not access status of transaction” failures.

3 http://wiki.postgresql.org/wiki/20110408pg_upgrade_fix
This error poses a significant risk of data loss for installations that have been upgraded with pg_upgrade. This patch corrects the problem for future uses of pg_upgrade, but does not in itself cure the issue in installations that have been processed with a buggy version of pg_upgrade.

- Suppress incorrect “PD_ALL_VISIBLE flag was incorrectly set” warning (Heikki Linnakangas)

VACUUM would sometimes issue this warning in cases that are actually valid.

- Use better SQLSTATE error codes for hot standby conflict cases (Tatsuo Ishii and Simon Riggs)

All retryable conflict errors now have an error code that indicates that a retry is possible. Also, session closure due to the database being dropped on the master is now reported as ERRCODE_DATABASE_DROPPED, rather than ERRCODE_ADMIN_SHUTDOWN, so that connection poolers can handle the situation correctly.

- Prevent intermittent hang in interactions of startup process with bgwriter process (Simon Riggs)

This affected recovery in non-hot-standby cases.

- Disallow including a composite type in itself (Tom Lane)

This prevents scenarios wherein the server could recurse infinitely while processing the composite type. While there are some possible uses for such a structure, they don't seem compelling enough to justify the effort required to make sure it always works safely.

- Avoid potential deadlock during catalog cache initialization (Nikhil Sontakke)

In some cases the cache loading code would acquire share lock on a system index before locking the index's catalog. This could deadlock against processes trying to acquire exclusive locks in the other, more standard order.

- Fix dangling-pointer problem in BEFORE ROW UPDATE trigger handling when there was a concurrent update to the target tuple (Tom Lane)

This bug has been observed to result in intermittent “cannot extract system attribute from virtual tuple” failures while trying to do UPDATE RETURNING ctid. There is a very small probability of more serious errors, such as generating incorrect index entries for the updated tuple.

- Disallow DROP TABLE when there are pending deferred trigger events for the table (Tom Lane)

Formerly the DROP would go through, leading to “could not open relation with OID nnn” errors when the triggers were eventually fired.

- Allow “replication” as a user name in pg_hba.conf (Andrew Dunstan)

“replication” is special in the database name column, but it was mistakenly also treated as special in the user name column.

- Prevent crash triggered by constant-false WHERE conditions during GEQO optimization (Tom Lane)

- Improve planner's handling of semi-join and anti-join cases (Tom Lane)

- Fix handling of SELECT FOR UPDATE in a sub-SELECT (Tom Lane)

This bug typically led to “cannot extract system attribute from virtual tuple” errors.

- Fix selectivity estimation for text search to account for NULLs (Jesper Krogh)

- Fix get_actual_variable_range() to support hypothetical indexes injected by an index adviser plugin (Gurjeet Singh)

2894
• Fix PL/Python memory leak involving array slices (Daniel Popowich)

• Allow libpq's SSL initialization to succeed when user's home directory is unavailable (Tom Lane)

If the SSL mode is such that a root certificate file is not required, there is no need to fail. This change restores the behavior to what it was in pre-9.0 releases.

• Fix libpq to return a useful error message for errors detected in conninfo_array_parse (Joseph Adams)

A typo caused the library to return NULL, rather than the PGconn structure containing the error message, to the application.

• Fix ecpg preprocessor's handling of float constants (Heikki Linnakangas)

• Fix parallel pg_restore to handle comments on POST_DATA items correctly (Arnd Hannemann)

• Fix pg_restore to cope with long lines (over 1KB) in TOC files (Tom Lane)

• Put in more safeguards against crashing due to division-by-zero with overly enthusiastic compiler optimization (Aurelien Jarno)

• Support use of dlopen() in FreeBSD and OpenBSD on MIPS (Tom Lane)

There was a hard-wired assumption that this system function was not available on MIPS hardware on these systems. Use a compile-time test instead, since more recent versions have it.

• Fix compilation failures on HP-UX (Heikki Linnakangas)

• Avoid crash when trying to write to the Windows console very early in process startup (Rushabh Lathia)

• Support building with MinGW 64 bit compiler for Windows (Andrew Dunstan)

• Fix version-incompatibility problem with libintl on Windows (Hiroshi Inoue)

• Fix usage of xcopy in Windows build scripts to work correctly under Windows 7 (Andrew Dunstan)

This affects the build scripts only, not installation or usage.

• Fix path separator used by pg_regress on Cygwin (Andrew Dunstan)

• Update time zone data files to tzdata release 2011f for DST law changes in Chile, Cuba, Falkland Islands, Morocco, Samoa, and Turkey; also historical corrections for South Australia, Alaska, and Hawaii.

E.155. Release 9.0.3

Release date: 2011-01-31

This release contains a variety of fixes from 9.0.2. For information about new features in the 9.0 major release, see Section E.158.

E.155.1. Migration to Version 9.0.3

A dump/restore is not required for those running 9.0.X.

E.155.2. Changes

• Before exiting walreceiver, ensure all the received WAL is fsync'd to disk (Heikki Linnakangas)
Otherwise the standby server could replay some un-synced WAL, conceivably leading to data corruption if the system crashes just at that point.

- Avoid excess fsync activity in walreceiver (Heikki Linnakangas)
- Make ALTER TABLE revalidate uniqueness and exclusion constraints when needed (Noah Misch)
  This was broken in 9.0 by a change that was intended to suppress revalidation during VACUUM FULL and CLUSTER, but unintentionally affected ALTER TABLE as well.
- Fix EvalPlanQual for UPDATE of an inheritance tree in which the tables are not all alike (Tom Lane)
  Any variation in the table row types (including dropped columns present in only some child tables) would confuse the EvalPlanQual code, leading to misbehavior or even crashes. Since EvalPlanQual is only executed during concurrent updates to the same row, the problem was only seen intermittently.
- Avoid failures when EXPLAIN tries to display a simple-form CASE expression (Tom Lane)
  If the CASE's test expression was a constant, the planner could simplify the CASE into a form that confused the expression-display code, resulting in “unexpected CASE WHEN clause” errors.
- Fix assignment to an array slice that is before the existing range of subscripts (Tom Lane)
  If there was a gap between the newly added subscripts and the first pre-existing subscript, the code miscalculated how many entries needed to be copied from the old array's null bitmap, potentially leading to data corruption or crash.
- Avoid unexpected conversion overflow in planner for very distant date values (Tom Lane)
  The date type supports a wider range of dates than can be represented by the timestamp types, but the planner assumed it could always convert a date to timestamp with impunity.
- Fix PL/Python crash when an array contains null entries (Alex Hunsaker)
- Remove ecpg's fixed length limit for constants defining an array dimension (Michael Meskes)
- Fix erroneous parsing of tsquery values containing ... & !(subexpression) | ...
  (Tom Lane)
  Queries containing this combination of operators were not executed correctly. The same error existed in contrib/intarray's query_int type and contrib/ltree's ltxquery type.
- Fix buffer overrun in contrib/intarray's input function for the query_int type (Apple)
  This bug is a security risk since the function's return address could be overwritten. Thanks to Apple Inc's security team for reporting this issue and supplying the fix. (CVE-2010-4015)
- Fix bug in contrib/seg's GiST picksplit algorithm (Alexander Korotkov)
  This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a seg column. If you have such an index, consider REINDEXing it after installing this update. (This is identical to the bug that was fixed in contrib/cube in the previous update.)

E.156. Release 9.0.2

Release date: 2010-12-16

This release contains a variety of fixes from 9.0.1. For information about new features in the 9.0 major release, see Section E.158.
E.156.1. Migration to Version 9.0.2

A dump/restore is not required for those running 9.0.X.

E.156.2. Changes

- Force the default `wal_sync_method` to be `fdatasync` on Linux (Tom Lane, Marti Raudsepp)

The default on Linux has actually been `fdatasync` for many years, but recent kernel changes caused PostgreSQL to choose `open_dataasync` instead. This choice did not result in any performance improvement, and caused outright failures on certain filesystems, notably `ext4` with the `data=journal` mount option.

- Fix “too many KnownAssignedXids” error during Hot Standby replay (Heikki Linnakangas)

- Fix race condition in lock acquisition during Hot Standby (Simon Riggs)

- Avoid unnecessary conflicts during Hot Standby (Simon Riggs)

  This fixes some cases where replay was considered to conflict with standby queries (causing delay of replay or possibly cancellation of the queries), but there was no real conflict.

- Fix assorted bugs in WAL replay logic for GIN indexes (Tom Lane)

  This could result in “bad buffer id: 0” failures or corruption of index contents during replication.

- Fix recovery from base backup when the starting checkpoint WAL record is not in the same WAL segment as its redo point (Jeff Davis)

- Fix corner-case bug when streaming replication is enabled immediately after creating the master database cluster (Heikki Linnakangas)

- Fix persistent slowdown of autovacuum workers when multiple workers remain active for a long time (Tom Lane)

  The effective `vacuum_cost_limit` for an autovacuum worker could drop to nearly zero if it processed enough tables, causing it to run extremely slowly.

- Fix long-term memory leak in autovacuum launcher (Alvaro Herrera)

- Avoid failure when trying to report an impending transaction wraparound condition from outside a transaction (Tom Lane)

  This oversight prevented recovery after transaction wraparound got too close, because database startup processing would fail.

- Add support for detecting register-stack overrun on IA64 (Tom Lane)

  The IA64 architecture has two hardware stacks. Full prevention of stack-overrun failures requires checking both.

- Add a check for stack overflow in `copyObject()` (Tom Lane)

  Certain code paths could crash due to stack overflow given a sufficiently complex query.

- Fix detection of page splits in temporary GiST indexes (Heikki Linnakangas)

  It is possible to have a “concurrent” page split in a temporary index, if for example there is an open cursor scanning the index when an insertion is done. GiST failed to detect this case and hence could deliver wrong results when execution of the cursor continued.

- Fix error checking during early connection processing (Tom Lane)
The check for too many child processes was skipped in some cases, possibly leading to postmaster crash when attempting to add the new child process to fixed-size arrays.

- Improve efficiency of window functions (Tom Lane)

  Certain cases where a large number of tuples needed to be read in advance, but work_mem was large enough to allow them all to be held in memory, were unexpectedly slow. percent_rank(), cume_dist() and ntile() in particular were subject to this problem.

- Avoid memory leakage while ANALYZE'ing complex index expressions (Tom Lane)

- Ensure an index that uses a whole-row Var still depends on its table (Tom Lane)

  An index declared like create index i on t (foo(t.*)) would not automatically get dropped when its table was dropped.

- Add missing support in DROP OWNED BY for removing foreign data wrapper/server privileges belonging to a user (Heikki Linnakangas)

- Do not “inline” a SQL function with multiple OUT parameters (Tom Lane)

  This avoids a possible crash due to loss of information about the expected result rowtype.

- Fix crash when inline-ing a set-returning function whose argument list contains a reference to an inline-able user function (Tom Lane)

- Behave correctly if ORDER BY, LIMIT, FOR UPDATE, or WITH is attached to the VALUES part of INSERT ... VALUES (Tom Lane)

- Make the OFF keyword unreserved (Heikki Linnakangas)

  This prevents problems with using off as a variable name in PL/pgSQL. That worked before 9.0, but was now broken because PL/pgSQL now treats all core reserved words as reserved.

- Fix constant-folding of COALESCE() expressions (Tom Lane)

  The planner would sometimes attempt to evaluate sub-expressions that in fact could never be reached, possibly leading to unexpected errors.

- Fix “could not find pathkey item to sort” planner failure with comparison of whole-row Vars (Tom Lane)

- Fix postmaster crash when connection acceptance (accept() or one of the calls made immediately after it) fails, and the postmaster was compiled with GSSAPI support (Alexander Chernikov)

- Retry after receiving an invalid response packet from a RADIUS authentication server (Magnus Hagander)

  This fixes a low-risk potential denial of service condition.

- Fix missed unlink of temporary files when log_temp_files is active (Tom Lane)

  If an error occurred while attempting to emit the log message, the unlink was not done, resulting in accumulation of temp files.

- Add print functionality for InhRelation nodes (Tom Lane)

  This avoids a failure when debug_print_parse is enabled and certain types of query are executed.

- Fix incorrect calculation of distance from a point to a horizontal line segment (Tom Lane)

  This bug affected several different geometric distance-measurement operators.
• Fix incorrect calculation of transaction status in ecpg (Itagaki Takahiro)
• Fix errors in psql's Unicode-escape support (Tom Lane)
• Speed up parallel pg_restore when the archive contains many large objects (blobs) (Tom Lane)
• Fix PL/pgSQL's handling of “simple” expressions to not fail in recursion or error-recovery cases (Tom Lane)
• Fix PL/pgSQL's error reporting for no-such-column cases (Tom Lane)
  As of 9.0, it would sometimes report “missing FROM-clause entry for table foo” when “record foo has no field bar” would be more appropriate.
• Fix PL/Python to honor typmod (i.e., length or precision restrictions) when assigning to tuple fields (Tom Lane)
  This fixes a regression from 8.4.
• Fix PL/Python's handling of set-returning functions (Jan Urbanski)
  Attempts to call SPI functions within the iterator generating a set result would fail.
• Fix bug in contrib/cube's GiST picksplit algorithm (Alexander Korotkov)
  This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a cube column. If you have such an index, consider REINDEXing it after installing this update.
• Don't emit “identifier will be truncated” notices in contrib/dblink except when creating new connections (Itagaki Takahiro)
• Fix potential coredump on missing public key in contrib/pgcrypto (Marti Raudsepp)
• Fix buffer overrun in contrib/pg_upgrade (Hernan Gonzalez)
• Fix memory leak in contrib/xml2's XPath query functions (Tom Lane)
• Update time zone data files to tzdata release 2010o for DST law changes in Fiji and Samoa; also historical corrections for Hong Kong.

E.157. Release 9.0.1

Release date: 2010-10-04

This release contains a variety of fixes from 9.0.0. For information about new features in the 9.0 major release, see Section E.158.

E.157.1. Migration to Version 9.0.1

A dump/restore is not required for those running 9.0.X.

E.157.2. Changes

• Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

  This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function's owner.
The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

- Improve `pg_get_expr()` security fix so that the function can still be used on the output of a sub-select (Tom Lane)
- Fix incorrect placement of placeholder evaluation (Tom Lane)
  This bug could result in query outputs being non-null when they should be null, in cases where the inner side of an outer join is a sub-select with non-strict expressions in its output list.
- Fix join removal's handling of placeholder expressions (Tom Lane)
- Fix possible duplicate scans of `UNION ALL` member relations (Tom Lane)
- Prevent infinite loop in `ProcessIncomingNotify()` after unlistening (Jeff Davis)
- Prevent `show_session_authorization()` from crashing within autovacuum processes (Tom Lane)
- Re-allow input of Julian dates prior to 0001-01-01 AD (Tom Lane)
  Input such as `'J100000':::date` worked before 8.4, but was unintentionally broken by added error-checking.
- Make `psql` recognize `DISCARD ALL` as a command that should not be encased in a transaction block in autocommit-off mode (Itagaki Takahiro)
- Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)

**E.158. Release 9.0**

**Release date:** 2010-09-20

**E.158.1. Overview**

This release of PostgreSQL adds features that have been requested for years, such as easy-to-use replication, a mass permission-changing facility, and anonymous code blocks. While past major releases have been conservative in their scope, this release shows a bold new desire to provide facilities that new and existing users of PostgreSQL will embrace. This has all been done with few incompatibilities. Major enhancements include:

- Built-in replication based on log shipping. This advance consists of two features: Streaming Replication, allowing continuous archive (WAL) files to be streamed over a network connection to a standby server, and Hot Standby, allowing continuous archive standby servers to execute read-only queries. The net effect is to support a single master with multiple read-only slave servers.
- Easier database object permissions management. `GRANT/REVOKE IN SCHEMA` supports mass permissions changes on existing objects, while `ALTER DEFAULT PRIVILEGES` allows control of privileges for objects created in the future. Large objects (BLOBs) now support permissions management as well.
• Broadly enhanced stored procedure support. The `DO` statement supports ad-hoc or “anonymous” code blocks. Functions can now be called using named parameters. PL/pgSQL is now installed by default, and PL/Perl and PL/Python have been enhanced in several ways, including support for Python3.

• Full support for 64-bit Windows.

• More advanced reporting queries, including additional windowing options (`PRECEDING` and `FOLLOWING`) and the ability to control the order in which values are fed to aggregate functions.

• New trigger features, including SQL-standard-compliant per-column triggers and conditional trigger execution.

• Deferrable unique constraints. Mass updates to unique keys are now possible without trickery.

• Exclusion constraints. These provide a generalized version of unique constraints, allowing enforcement of complex conditions.

• New and enhanced security features, including RADIUS authentication, LDAP authentication improvements, and a new contrib module `passwordcheck` for testing password strength.

• New high-performance implementation of the `LISTEN`/`NOTIFY` feature. Pending events are now stored in a memory-based queue rather than a table. Also, a “payload” string can be sent with each event, rather than transmitting just an event name as before.

• New implementation of `VACUUM FULL`. This command now rewrites the entire table and indexes, rather than moving individual rows to compact space. It is substantially faster in most cases, and no longer results in index bloat.

• New contrib module `pg_upgrade` to support in-place upgrades from 8.3 or 8.4 to 9.0.

• Multiple performance enhancements for specific types of queries, including elimination of unnecessary joins. This helps optimize some automatically-generated queries, such as those produced by object-relational mappers (ORMs).

• `EXPLAIN` enhancements. The output is now available in JSON, XML, or YAML format, and includes buffer utilization and other data not previously available.

• `hstore` improvements, including new functions and greater data capacity.

The above items are explained in more detail in the sections below.

E.158.2. Migration to Version 9.0

A dump/restore using `pg_dump`, or use of `pg_upgrade`, is required for those wishing to migrate data from any previous release.

Version 9.0 contains a number of changes that selectively break backwards compatibility in order to support new features and code quality improvements. In particular, users who make extensive use of PL/pgSQL, Point-In-Time Recovery (PITR), or Warm Standby should test their applications because of slight user-visible changes in those areas. Observe the following incompatibilities:

E.158.2.1. Server Settings

• Remove server parameter `add_missing_from`, which was defaulted to off for many years (Tom Lane)

• Remove server parameter `regex_flavor`, which was defaulted to `advanced` for many years (Tom Lane)

• `archive_mode` now only affects `archive_command`; a new setting, `wal_level`, affects the contents of the write-ahead log (Heikki Linnakangas)
Release Notes

• log_temp_files now uses default file size units of kilobytes (Robert Haas)

E.158.2.2. Queries

• When querying a parent table, do not do any separate permission checks on child tables scanned as part of the query (Peter Eisentraut)

The SQL standard specifies this behavior, and it is also much more convenient in practice than the former behavior of checking permissions on each child as well as the parent.

E.158.2.3. Data Types

• bytea output now appears in hex format by default (Peter Eisentraut)

The server parameter bytea_output can be used to select the traditional output format if needed for compatibility.

• Array input now considers only plain ASCII whitespace characters to be potentially ignorable; it will never ignore non-ASCII characters, even if they are whitespace according to some locales (Tom Lane)

This avoids some corner cases where array values could be interpreted differently depending on the server's locale settings.

• Improve standards compliance of SIMILAR TO patterns and SQL-style substring() patterns (Tom Lane)

This includes treating ? and {...} as pattern metacharacters, while they were simple literal characters before; that corresponds to new features added in SQL:2008. Also, ^ and $ are now treated as simple literal characters; formerly they were treated as metacharacters, as if the pattern were following POSIX rather than SQL rules. Also, in SQL-standard substring(), use of parentheses for nesting no longer interferes with capturing of a substring. Also, processing of bracket expressions (character classes) is now more standards-compliant.

• Reject negative length values in 3-parameter substring() for bit strings, per the SQL standard (Tom Lane)

• Make date_trunc truncate rather than round when reducing precision of fractional seconds (Tom Lane)

The code always acted this way for integer-based dates/times. Now float-based dates/times behave similarly.

E.158.2.4. Object Renaming

• Tighten enforcement of column name consistency during RENAME when a child table inherits the same column from multiple unrelated parents (KaiGai Kohei)

• No longer automatically rename indexes and index columns when the underlying table columns are renamed (Tom Lane)

Administrators can still rename such indexes and columns manually. This change will require an update of the JDBC driver, and possibly other drivers, so that unique indexes are correctly recognized after a rename.

• CREATE OR REPLACE FUNCTION can no longer change the declared names of function parameters (Pavel Stehule)

In order to avoid creating ambiguity in named-parameter calls, it is no longer allowed to change the aliases for input parameters in the declaration of an existing function (although names can still
be assigned to previously unnamed parameters). You now have to DROP and recreate the function to do that.

E.158.2.5. PL/pgSQL

• PL/pgSQL now throws an error if a variable name conflicts with a column name used in a query (Tom Lane)

The former behavior was to bind ambiguous names to PL/pgSQL variables in preference to query columns, which often resulted in surprising misbehavior. Throwing an error allows easy detection of ambiguous situations. Although it's recommended that functions encountering this type of error be modified to remove the conflict, the old behavior can be restored if necessary via the configuration parameter plpgsql.variable_conflict, or via the per-function option #variable_conflict.

• PL/pgSQL no longer allows variable names that match certain SQL reserved words (Tom Lane)

This is a consequence of aligning the PL/pgSQL parser to match the core SQL parser more closely. If necessary, variable names can be double-quoted to avoid this restriction.

• PL/pgSQL now requires columns of composite results to match the expected type modifier as well as base type (Pavel Stehule, Tom Lane)

For example, if a column of the result type is declared as \texttt{NUMERIC(30,2)}, it is no longer acceptable to return a \texttt{NUMERIC} of some other precision in that column. Previous versions neglected to check the type modifier and would thus allow result rows that didn't actually conform to the declared restrictions.

• PL/pgSQL now treats selection into composite fields more consistently (Tom Lane)

Formerly, a statement like \texttt{SELECT ... INTO rec.fld FROM ...} was treated as a scalar assignment even if the record field \texttt{fld} was of composite type. Now it is treated as a record assignment, the same as when the INTO target is a regular variable of composite type. So the values to be assigned to the field's subfields should be written as separate columns of the SELECT list, not as a \texttt{ROW(...)} construct as in previous versions.

If you need to do this in a way that will work in both 9.0 and previous releases, you can write something like \texttt{rec.fld ::= ROW(...) FROM ....}

• Remove PL/pgSQL's RENAME declaration (Tom Lane)

Instead of RENAME, use ALIAS, which can now create an alias for any variable, not only dollar sign parameter names (such as \$1) as before.

E.158.2.6. Other Incompatibilities

• Deprecate use of => as an operator name (Robert Haas)

Future versions of PostgreSQL will probably reject this operator name entirely, in order to support the SQL-standard notation for named function parameters. For the moment, it is still allowed, but a warning is emitted when such an operator is defined.

• Remove support for platforms that don't have a working 64-bit integer data type (Tom Lane)

It is believed all still-supported platforms have working 64-bit integer data types.

E.158.3. Changes

Version 9.0 has an unprecedented number of new major features, and over 200 enhancements, improvements, new commands, new functions, and other changes.
E.158.3.1. Server

E.158.3.1.1. Continuous Archiving and Streaming Replication

PostgreSQL's existing standby-server capability has been expanded both to support read-only queries on standby servers and to greatly reduce the lag between master and standby servers. For many users, this will be a useful and low-administration form of replication, either for high availability or for horizontal scalability.

- Allow a standby server to accept read-only queries (Simon Riggs, Heikki Linnakangas)

  This feature is called Hot Standby. There are new `postgresql.conf` and `recovery.conf` settings to control this feature, as well as extensive documentation.

- Allow write-ahead log (WAL) data to be streamed to a standby server (Fujii Masao, Heikki Linnakangas)

  This feature is called Streaming Replication. Previously WAL data could be sent to standby servers only in units of entire WAL files (normally 16 megabytes each). Streaming Replication eliminates this inefficiency and allows updates on the master to be propagated to standby servers with very little delay. There are new `postgresql.conf` and `recovery.conf` settings to control this feature, as well as extensive documentation.

- Add `pg_last_xlog_receive_location()` and `pg_last_xlog_replay_location()`, which can be used to monitor standby server WAL activity (Simon Riggs, Fujii Masao, Heikki Linnakangas)

E.158.3.1.2. Performance

- Allow per-tablespace values to be set for sequential and random page cost estimates (`seq_page_cost/random_page_cost`) via `ALTER TABLESPACE ... SET/RESET` (Robert Haas)

- Improve performance and reliability of EvalPlanQual rechecks in join queries (Tom Lane)

  `UPDATE`, `DELETE`, and `SELECT FOR UPDATE/SHARE` queries that involve joins will now behave much better when encountering freshly-updated rows.

- Improve performance of `TRUNCATE` when the table was created or truncated earlier in the same transaction (Tom Lane)

- Improve performance of finding inheritance child tables (Tom Lane)

E.158.3.1.3. Optimizer

- Remove unnecessary outer joins (Robert Haas)

  Outer joins where the inner side is unique and not referenced above the join are unnecessary and are therefore now removed. This will accelerate many automatically generated queries, such as those created by object-relational mappers (ORMs).

- Allow `IS NOT NULL` restrictions to use indexes (Tom Lane)

  This is particularly useful for finding `MAX() / MIN()` values in indexes that contain many null values.

- Improve the optimizer's choices about when to use materialize nodes, and when to use sorting versus hashing for `DISTINCT` (Tom Lane)

- Improve the optimizer's equivalence detection for expressions involving `boolean <>` operators (Tom Lane)
E.158.3.1.4. GEQO

- Use the same random seed every time GEQO plans a query (Andres Freund)

  While the Genetic Query Optimizer (GEQO) still selects random plans, it now always selects the
  same random plans for identical queries, thus giving more consistent performance. You can modify
  geqo_seed to experiment with alternative plans.

- Improve GEQO plan selection (Tom Lane)

  This avoids the rare error “failed to make a valid plan”, and should also improve planning speed.

E.158.3.1.5. Optimizer Statistics

- Improve ANALYZE to support inheritance-tree statistics (Tom Lane)

  This is particularly useful for partitioned tables. However, autovacuum does not yet automatically
  re-analyze parent tables when child tables change.

- Improve autovacuum’s detection of when re-analyze is necessary (Tom Lane)

- Improve optimizer’s estimation for greater/less-than comparisons (Tom Lane)

  When looking up statistics for greater/less-than comparisons, if the comparison value is in the first
  or last histogram bucket, use an index (if available) to fetch the current actual column minimum or
  maximum. This greatly improves the accuracy of estimates for comparison values near the ends of
  the data range, particularly if the range is constantly changing due to addition of new data.

- Allow setting of number-of-distinct-values statistics using ALTER TABLE (Robert Haas)

  This allows users to override the estimated number or percentage of distinct values for a column.
  This statistic is normally computed by ANALYZE, but the estimate can be poor, especially on tables
  with very large numbers of rows.

E.158.3.1.6. Authentication

- Add support for RADIUS (Remote Authentication Dial In User Service) authentication (Magnus
  Hagander)

- Allow LDAP (Lightweight Directory Access Protocol) authentication to operate in “search/bind”
  mode (Robert Fleming, Magnus Hagander)

  This allows the user to be looked up first, then the system uses the DN (Distinguished Name) re-
  turned for that user.

- Add samehost and samenet designations to pg_hba.conf (Stef Walter)

  These match the server’s IP address and subnet address respectively.

- Pass trusted SSL root certificate names to the client so the client can return an appropriate client
  certificate (Craig Ringer)

E.158.3.1.7. Monitoring

- Add the ability for clients to set an application name, which is displayed in pg_stat_activity
  (Dave Page)

  This allows administrators to characterize database traffic and troubleshoot problems by source
  application.

- Add a SQLSTATE option (%e) to log_line_prefix (Guillaume Smet)
This allows users to compile statistics on errors and messages by error code number.

- Write to the Windows event log in UTF16 encoding (Itagaki Takahiro)

Now there is true multilingual support for PostgreSQL log messages on Windows.

### E.158.3.1.8. Statistics Counters

- Add `pg_stat_reset_shared('bgwriter')` to reset the cluster-wide shared statistics for the background writer (Greg Smith)

- Add `pg_stat_reset_single_table_counters()` and `pg_stat_reset_single_function_counters()` to allow resetting the statistics counters for individual tables and functions (Magnus Hagander)

### E.158.3.1.9. Server Settings

- Allow setting of configuration parameters based on database/role combinations (Alvaro Herrera)

  Previously only per-database and per-role settings were possible, not combinations. All role and database settings are now stored in the new `pg_db_role_setting` system catalog. A new `psql` command `\drds` shows these settings. The legacy system views `pg_roles`, `pg_shadow`, and `pg_user` do not show combination settings, and therefore no longer completely represent the configuration for a user or database.

- Add server parameter `bonjour`, which controls whether a Bonjour-enabled server advertises itself via Bonjour (Tom Lane)

  The default is off, meaning it does not advertise. This allows packagers to distribute Bonjour-enabled builds without worrying that individual users might not want the feature.

- Add server parameter `enable_material`, which controls the use of materialize nodes in the optimizer (Robert Haas)

  The default is on. When off, the optimizer will not add materialize nodes purely for performance reasons, though they will still be used when necessary for correctness.

- Change server parameter `log_temp_files` to use default file size units of kilobytes (Robert Haas)

  Previously this setting was interpreted in bytes if no units were specified.

- Log changes of parameter values when `postgresql.conf` is reloaded (Peter Eisentraut)

  This lets administrators and security staff audit changes of database settings, and is also very convenient for checking the effects of `postgresql.conf` edits.

- Properly enforce superuser permissions for custom server parameters (Tom Lane)

  Non-superusers can no longer issue `ALTER ROLE/DATABASE SET` for parameters that are not currently known to the server. This allows the server to correctly check that superuser-only parameters are only set by superusers. Previously, the `SET` would be allowed and then ignored at session start, making superuser-only custom parameters much less useful than they should be.

### E.158.3.2. Queries

- Perform `SELECT FOR UPDATE/SHARE` processing after applying `LIMIT`, so the number of rows returned is always predictable (Tom Lane)

  Previously, changes made by concurrent transactions could cause a `SELECT FOR UPDATE` to unexpectedly return fewer rows than specified by its `LIMIT`. `FOR UPDATE` in combination with
ORDER BY can still produce surprising results, but that can be corrected by placing FOR UPDATE in a subquery.

- Allow mixing of traditional and SQL-standard LIMIT/OFFSET syntax (Tom Lane)
- Extend the supported frame options in window functions (Hitoshi Harada)

Frames can now start with CURRENT ROW, and the ROWS n PRECEDING/FOLLOWING options are now supported.

- Make SELECT INTO and CREATE TABLE AS return row counts to the client in their command tags (Boszormenyi Zoltan)

This can save an entire round-trip to the client, allowing result counts and pagination to be calculated without an additional COUNT query.

### E.158.3.2.1. Unicode Strings

- Support Unicode surrogate pairs (dual 16-bit representation) in U& strings and identifiers (Peter Eisentraut)

- Support Unicode escapes in E'...' strings (Marko Kreen)

### E.158.3.3. Object Manipulation

- Speed up CREATE DATABASE by deferring flushes to disk (Andres Freund, Greg Stark)

- Allow comments on columns of tables, views, and composite types only, not other relation types such as indexes and TOAST tables (Tom Lane)

- Allow the creation of enumerated types containing no values (Bruce Momjian)

- Let values of columns having storage type MAIN remain on the main heap page unless the row cannot fit on a page (Kevin Grittner)

Previously MAIN values were forced out to TOAST tables until the row size was less than one-quarter of the page size.

#### E.158.3.3.1. ALTER TABLE

- Implement IF EXISTS for ALTER TABLE DROP COLUMN and ALTER TABLE DROP CONSTRAINT (Andres Freund)

- Allow ALTER TABLE commands that rewrite tables to skip WAL logging (Itagaki Takahiro)

Such operations either produce a new copy of the table or are rolled back, so WAL archiving can be skipped, unless running in continuous archiving mode. This reduces I/O overhead and improves performance.

- Fix failure of ALTER TABLE table ADD COLUMN col serial when done by non-owner of table (Tom Lane)

#### E.158.3.3.2. CREATE TABLE

- Add support for copying COMMENTS and STORAGE settings in CREATE TABLE ... LIKE commands (Itagaki Takahiro)

- Add a shortcut for copying all properties in CREATE TABLE ... LIKE commands (Itagaki Takahiro)

- Add the SQL-standard CREATE TABLE ... OF type command (Peter Eisentraut)
This allows creation of a table that matches an existing composite type. Additional constraints and defaults can be specified in the command.

### E.158.3.3.3. Constraints

- **Add deferrable unique constraints (Dean Rasheed)**

  This allows mass updates, such as UPDATE tab SET col = col + 1, to work reliably on columns that have unique indexes or are marked as primary keys. If the constraint is specified as DEFERRABLE it will be checked at the end of the statement, rather than after each row is updated. The constraint check can also be deferred until the end of the current transaction, allowing such updates to be spread over multiple SQL commands.

- **Add exclusion constraints (Jeff Davis)**

  Exclusion constraints generalize uniqueness constraints by allowing arbitrary comparison operators, not just equality. They are created with the CREATE TABLE CONSTRAINT ... EXCLUDE clause. The most common use of exclusion constraints is to specify that column entries must not overlap, rather than simply not be equal. This is useful for time periods and other ranges, as well as arrays. This feature enhances checking of data integrity for many calendaring, time-management, and scientific applications.

- **Improve uniqueness-constraint violation error messages to report the values causing the failure (Itagaki Takahiro)**

  For example, a uniqueness constraint violation might now report Key (x) = (2) already exists.

### E.158.3.3.4. Object Permissions

- **Add the ability to make mass permission changes across a whole schema using the new GRANT/REVOKE IN SCHEMA clause (Petr Jelinek)**

  This simplifies management of object permissions and makes it easier to utilize database roles for application data security.

- **Add ALTER DEFAULT PRIVILEGES command to control privileges of objects created later (Petr Jelinek)**

  This greatly simplifies the assignment of object privileges in a complex database application. Default privileges can be set for tables, views, sequences, and functions. Defaults may be assigned on a per-schema basis, or database-wide.

- **Add the ability to control large object (BLOB) permissions with GRANT/REVOKE (KaiGai Kohei)**

  Formerly, any database user could read or modify any large object. Read and write permissions can now be granted and revoked per large object, and the ownership of large objects is tracked.

### E.158.3.4. Utility Operations

- **Make LISTEN/NOTIFY store pending events in a memory queue, rather than in a system table (Joachim Wieland)**

  This substantially improves performance, while retaining the existing features of transactional support and guaranteed delivery.

- **Allow NOTIFY to pass an optional “payload” string to listeners (Joachim Wieland)**

  This greatly improves the usefulness of LISTEN/NOTIFY as a general-purpose event queue system.
• Allow CLUSTER on all per-database system catalogs (Tom Lane)

Shared catalogs still cannot be clustered.

**E.158.3.4.1. COPY**

• Accept COPY ... CSV FORCE QUOTE * (Itagaki Takahiro)

Now * can be used as shorthand for “all columns” in the FORCE QUOTE clause.

• Add new COPY syntax that allows options to be specified inside parentheses (Robert Haas, Emmanuel Cecchet)

This allows greater flexibility for future COPY options. The old syntax is still supported, but only for pre-existing options.

**E.158.3.4.2. EXPLAIN**

• Allow EXPLAIN to output in XML, JSON, or YAML format (Robert Haas, Greg Sabino Mullane)

The new output formats are easily machine-readable, supporting the development of new tools for analysis of EXPLAIN output.

• Add new BUFFERS option to report query buffer usage during EXPLAIN ANALYZE (Itagaki Takahiro)

This allows better query profiling for individual queries. Buffer usage is no longer reported in the output for log_statement_stats and related settings.

• Add hash usage information to EXPLAIN output (Robert Haas)

• Add new EXPLAIN syntax that allows options to be specified inside parentheses (Robert Haas)

This allows greater flexibility for future EXPLAIN options. The old syntax is still supported, but only for pre-existing options.

**E.158.3.4.3. VACUUM**

• Change VACUUM FULL to rewrite the entire table and rebuild its indexes, rather than moving individual rows around to compact space (Itagaki Takahiro, Tom Lane)

The previous method was usually slower and caused index bloat. Note that the new method will use more disk space transiently during VACUUM FULL; potentially as much as twice the space normally occupied by the table and its indexes.

• Add new VACUUM syntax that allows options to be specified inside parentheses (Itagaki Takahiro)

This allows greater flexibility for future VACUUM options. The old syntax is still supported, but only for pre-existing options.

**E.158.3.4.4. Indexes**

• Allow an index to be named automatically by omitting the index name in CREATE INDEX (Tom Lane)

• By default, multicolumn indexes are now named after all their columns; and index expression columns are now named based on their expressions (Tom Lane)

• Reindexing shared system catalogs is now fully transactional and crash-safe (Tom Lane)

Formerly, reindexing a shared index was only allowed in standalone mode, and a crash during the operation could leave the index in worse condition than it was before.
• Add `point_ops` operator class for GiST (Teodor Sigaev)

This feature permits GiST indexing of `point` columns. The index can be used for several types of queries such as `point <@ polygon` (point is in polygon). This should make many PostGIS queries faster.

• Use red-black binary trees for GIN index creation (Teodor Sigaev)

Red-black trees are self-balancing. This avoids slowdowns in cases where the input is in nonrandom order.

E.158.3.5. Data Types

• Allow `bytea` values to be written in hex notation (Peter Eisentraut)

The server parameter `bytea_output` controls whether hex or traditional format is used for `bytea` output. Libpq's `PQescapeByteaConn()` function automatically uses the hex format when connected to PostgreSQL 9.0 or newer servers. However, pre-9.0 libpq versions will not correctly process hex format from newer servers.

The new hex format will be directly compatible with more applications that use binary data, allowing them to store and retrieve it without extra conversion. It is also significantly faster to read and write than the traditional format.

• Allow server parameter `extra_float_digits` to be increased to 3 (Tom Lane)

The previous maximum `extra_float_digits` setting was 2. There are cases where 3 digits are needed to dump and restore `float4` values exactly. `pg_dump` will now use the setting of 3 when dumping from a server that allows it.

• Tighten input checking for `int2vector` values (Caleb Welton)

E.158.3.5.1. Full Text Search

• Add prefix support in `synonym` dictionaries (Teodor Sigaev)

• Add `filtering` dictionaries (Teodor Sigaev)

  Filtering dictionaries allow tokens to be modified then passed to subsequent dictionaries.

• Allow underscores in email-address tokens (Teodor Sigaev)

• Use more standards-compliant rules for parsing URL tokens (Tom Lane)

E.158.3.6. Functions

• Allow function calls to supply parameter names and match them to named parameters in the function definition (Pavel Stehule)

  For example, if a function is defined to take parameters `a` and `b`, it can be called with `func(a := 7, b := 12)` or `func(b := 12, a := 7)`.

• Support locale-specific regular expression processing with UTF-8 server encoding (Tom Lane)

  Locale-specific regular expression functionality includes case-insensitive matching and locale-specific character classes. Previously, these features worked correctly for non-ASCII characters only if the database used a single-byte server encoding (such as LATIN1). They will still misbehave in multi-byte encodings other than UTF-8.

• Add support for scientific notation in `to_char()` (EEE specification) (Pavel Stehule, Brendan Jurd)
• Make `to_char()` honor FM (fill mode) in Y, YY, and YYY specifications (Bruce Momjian, Tom Lane)

It was already honored by YYYY.

• Fix `to_char()` to output localized numeric and monetary strings in the correct encoding on Windows (Hiroshi Inoue, Itagaki Takahiro, Bruce Momjian)

• Correct calculations of “overlaps” and “contains” operations for polygons (Teodor Sigaev)

The polygon `&&` (overlaps) operator formerly just checked to see if the two polygons' bounding boxes overlapped. It now does a more correct check. The polygon `@>` and `<@` (contains/contained by) operators formerly checked to see if one polygon's vertexes were all contained in the other; this can wrongly report “true” for some non-convex polygons. Now they check that all line segments of one polygon are contained in the other.

E.158.3.6.1. Aggregates

• Allow aggregate functions to use `ORDER BY` (Andrew Gierth)

For example, this is now supported: `array_agg(a ORDER BY b)`. This is useful with aggregates for which the order of input values is significant, and eliminates the need to use a nonstandard subquery to determine the ordering.

• Multi-argument aggregate functions can now use `DISTINCT` (Andrew Gierth)

• Add the `string_agg()` aggregate function to combine values into a single string (Pavel Stehule)

• Aggregate functions that are called with `DISTINCT` are now passed NULL values if the aggregate transition function is not marked as `STRICT` (Andrew Gierth)

For example, `agg(DISTINCT x)` might pass a NULL `x` value to `agg()`. This is more consistent with the behavior in non-`DISTINCT` cases.

E.158.3.6.2. Bit Strings

• Add `get_bit()` and `set_bit()` functions for bit strings, mirroring those for `bytea` (Leonardo F)

• Implement `OVERLAY()` (replace) for bit strings and `bytea` (Leonardo F)

E.158.3.6.3. Object Information Functions

• Add `pg_table_size()` and `pg_indexes_size()` to provide a more user-friendly interface to the `pg_relation_size()` function (Bernd Helmle)

• Add `has_sequence_privilege()` for sequence permission checking (Abhijit Menon-Sen)

• Update the `information_schema` views to conform to SQL:2008 (Peter Eisenraut)

• Make the `information_schema` views correctly display maximum octet lengths for `char` and `varchar` columns (Peter Eisenraut)

• Speed up `information_schema` privilege views (Joachim Wieland)

E.158.3.6.4. Function and Trigger Creation

• Support execution of anonymous code blocks using the `DO` statement (Petr Jelinek, Joshua Tolley, Hannu Valtonen)

This allows execution of server-side code without the need to create and delete a temporary function definition. Code can be executed in any language for which the user has permissions to define a function.
• Implement SQL-standard-compliant per-column triggers (Itagaki Takahiro)

Such triggers are fired only when the specified column(s) are affected by the query, e.g. appear in an UPDATE's SET list.

• Add the WHEN clause to CREATE TRIGGER to allow control over whether a trigger is fired (Itagaki Takahiro)

While the same type of check can always be performed inside the trigger, doing it in an external WHEN clause can have performance benefits.

E.158.3.7. Server-Side Languages

• Add the OR REPLACE clause to CREATE LANGUAGE (Tom Lane)

This is helpful to optionally install a language if it does not already exist, and is particularly helpful now that PL/pgSQL is installed by default.

E.158.3.7.1. PL/pgSQL Server-Side Language

• Install PL/pgSQL by default (Bruce Momjian)

The language can still be removed from a particular database if the administrator has security or performance concerns about making it available.

• Improve handling of cases where PL/pgSQL variable names conflict with identifiers used in queries within a function (Tom Lane)

The default behavior is now to throw an error when there is a conflict, so as to avoid surprising behaviors. This can be modified, via the configuration parameter plpgsql.variable_conflict or the per-function option #variable_conflict, to allow either the variable or the query-supplied column to be used. In any case PL/pgSQL will no longer attempt to substitute variables in places where they would not be syntactically valid.

• Make PL/pgSQL use the main lexer, rather than its own version (Tom Lane)

This ensures accurate tracking of the main system's behavior for details such as string escaping. Some user-visible details, such as the set of keywords considered reserved in PL/pgSQL, have changed in consequence.

• Avoid throwing an unnecessary error for an invalid record reference (Tom Lane)

An error is now thrown only if the reference is actually fetched, rather than whenever the enclosing expression is reached. For example, many people have tried to do this in triggers:

if TG_OP = 'INSERT' and NEW.col1 = ... then

This will now actually work as expected.

• Improve PL/pgSQL's ability to handle row types with dropped columns (Pavel Stehule)

• Allow input parameters to be assigned values within PL/pgSQL functions (Steve Prentice)

Formerly, input parameters were treated as being declared CONST, so the function's code could not change their values. This restriction has been removed to simplify porting of functions from other DBMSes that do not impose the equivalent restriction. An input parameter now acts like a local variable initialized to the passed-in value.

• Improve error location reporting in PL/pgSQL (Tom Lane)

• Add count and ALL options to MOVE FORWARD/BACKWARD in PL/pgSQL (Pavel Stehule)
Release Notes

- Allow PL/pgSQL's `WHERE CURRENT OF` to use a cursor variable (Tom Lane)
- Allow PL/pgSQL's `OPEN cursor FOR EXECUTE` to use parameters (Pavel Stehule, Itagaki Takahiro)

This is accomplished with a new `USING` clause.

### E.158.3.7.2. PL/Perl Server-Side Language

- Add new PL/Perl functions: `quote_literal()`, `quote_nullable()`, `quote_identifier()`, `encode_bytea()`, `decode_bytea()`, `looks_like_number()`, `encode_array_literal()`, `encode_array_constructor()` (Tim Bunce)
- Add server parameter `plperl.on_init` to specify a PL/Perl initialization function (Tim Bunce)
  - `plperl.on_plperl_init` and `plperl.on_plperlu_init` are also available for initialization that is specific to the trusted or untrusted language respectively.
- Support `END` blocks in PL/Perl (Tim Bunce)
  - `END` blocks do not currently allow database access.
- Allow `use strict` in PL/Perl (Tim Bunce)
  - `Perl strict` checks can also be globally enabled with the new server parameter `plperl.use_strict`.
- Allow `require` in PL/Perl (Tim Bunce)
  - This basically tests to see if the module is loaded, and if not, generates an error. It will not allow loading of modules that the administrator has not preloaded via the initialization parameters.
- Allow `use feature` in PL/Perl if Perl version 5.10 or later is used (Tim Bunce)
- Verify that PL/Perl return values are valid in the server encoding (Andrew Dunstan)

### E.158.3.7.3. PL/Python Server-Side Language

- Add Unicode support in PL/Python (Peter Eisentraut)
  - Strings are automatically converted from/to the server encoding as necessary.
- Improve `bytea` support in PL/Python (Caleb Welton)
  - `Bytea` values passed into PL/Python are now represented as binary, rather than the PostgreSQL `bytea` text format. `Bytea` values containing null bytes are now also output properly from PL/Python. Passing of boolean, integer, and float values was also improved.
- Support arrays as parameters and return values in PL/Python (Peter Eisentraut)
- Improve mapping of SQL domains to Python types (Peter Eisentraut)
- Add Python 3 support to PL/Python (Peter Eisentraut)
  - The new server-side language is called `plpython3u`. This cannot be used in the same session with the Python 2 server-side language.
- Improve error location and exception reporting in PL/Python (Peter Eisentraut)

### E.158.3.8. Client Applications

- Add an `--analyze-only` option to `vacuumdb`, to analyze without vacuuming (Bruce Momjian)
E.158.3.8.1. psql

- Add support for quoting/escaping the values of psql variables as SQL strings or identifiers (Pavel Stehule, Robert Haas)

  For example, `:var` will produce the value of `var` quoted and properly escaped as a literal string, while `:"var"` will produce its value quoted and escaped as an identifier.

- Ignore a leading UTF-8-encoded Unicode byte-order marker in script files read by psql (Itagaki Takahiro)

  This is enabled when the client encoding is UTF-8. It improves compatibility with certain editors, mostly on Windows, that insist on inserting such markers.

- Fix `psql --file` to properly honor `--single-transaction` (Bruce Momjian)

- Avoid overwriting of psql's command-line history when two psql sessions are run concurrently (Tom Lane)

- Improve psql's tab completion support (Itagaki Takahiro)

- Show `\timing` output when it is enabled, regardless of “quiet” mode (Peter Eisentraut)

E.158.3.8.1.1. psql Display

- Improve display of wrapped columns in psql (Roger Leigh)

  This behavior is now the default. The previous formatting is available by using `\pset linestyle old-ascii`.

- Allow psql to use fancy Unicode line-drawing characters via `\pset linestyle unicode` (Roger Leigh)

E.158.3.8.1.2. psql `\d` Commands

- Make `\d` show child tables that inherit from the specified parent (Damien Clochard)

  `\d` shows only the number of child tables, while `\d+` shows the names of all child tables.

- Show definitions of index columns in `\d index_name` (Khee Chin)

  The definition is useful for expression indexes.

- Show a view’s defining query only in `\d+`, not in `\d` (Peter Eisentraut)

  Always including the query was deemed overly verbose.

E.158.3.8.2. pg_dump

- Make `pg_dump/pg_restore --clean` also remove large objects (Itagaki Takahiro)

- Fix `pg_dump` to properly dump large objects when `standard_conforming_strings` is enabled (Tom Lane)

  The previous coding could fail when dumping to an archive file and then generating script output from `pg_restore`.

- `pg_restore` now emits large-object data in hex format when generating script output (Tom Lane)

  This could cause compatibility problems if the script is then loaded into a pre-9.0 server. To work around that, restore directly to the server, instead.
• Allow pg_dump to dump comments attached to columns of composite types (Taro Minowa (Higepon))

• Make pg_dump --verbose output the pg_dump and server versions in text output mode (Jim Cox, Tom Lane)

These were already provided in custom output mode.

• pg_restore now complains if any command-line arguments remain after the switches and optional file name (Tom Lane)

Previously, it silently ignored any such arguments.

E.158.3.8.3. pg_ctl

• Allow pg_ctl to be used safely to start the postmaster during a system reboot (Tom Lane)

Previously, pg_ctl's parent process could have been mistakenly identified as a running postmaster based on a stale postmaster lock file, resulting in a transient failure to start the database.

• Give pg_ctl the ability to initialize the database (by invoking initdb) (Zdenek Kotala)

E.158.3.9. Development Tools

E.158.3.9.1. libpq

• Add new libpq functions PQconnectdbParams() and PQconnectStartParams() (Guillaume Lelarge)

These functions are similar to PQconnectdb() and PQconnectStart() except that they accept a null-terminated array of connection options, rather than requiring all options to be provided in a single string.

• Add libpq functions PQescapeLiteral() and PQescapeIdentifier() (Robert Haas)

These functions return appropriately quoted and escaped SQL string literals and identifiers. The caller is not required to pre-allocate the string result, as is required by PQescapeStringConn().

• Add support for a per-user service file (.pg_service.conf), which is checked before the site-wide service file (Peter Eisentraut)

• Properly report an error if the specified libpq service cannot be found (Peter Eisentraut)

• Add TCP keepalive settings in libpq (Tollef Fog Heen, Fujii Masao, Robert Haas)

Keepalive settings were already supported on the server end of TCP connections.

• Avoid extra system calls to block and unblock SIGPIPE in libpq, on platforms that offer alternative methods (Jeremy Kerr)

• When a .pgpass-supplied password fails, mention where the password came from in the error message (Bruce Momjian)

• Load all SSL certificates given in the client certificate file (Tom Lane)

This improves support for indirectly-signed SSL certificates.

E.158.3.9.2. ecpg

• Add SQLDA (SQL Descriptor Area) support to ecpg (Boszormenyi Zoltan)
• Add the `DESCRIBE [ OUTPUT ]` statement to `ecpg` (Boszormenyi Zoltan)

• Add an ECPGTransactionStatus function to return the current transaction status (Bernd Helmle)

• Add the `string` data type in `ecpg` Informix-compatibility mode (Boszormenyi Zoltan)

• Allow `ecpg` to use `new` and `old` variable names without restriction (Michael Meskes)

• Allow `ecpg` to use variable names in `free()` (Michael Meskes)

• Make `ecpg_dynamic_type()` return zero for non-SQL3 data types (Michael Meskes)

  Previously it returned the negative of the data type OID. This could be confused with valid type OIDs, however.

• Support `long long` types on platforms that already have 64-bit `long` (Michael Meskes)

E.158.3.9.2.1. `ecpg` Cursors

• Add out-of-scope cursor support in `ecpg`'s native mode (Boszormenyi Zoltan)

  This allows `DECLARE` to use variables that are not in scope when `OPEN` is called. This facility already existed in `ecpg`'s Informix-compatibility mode.

• Allow dynamic cursor names in `ecpg` (Boszormenyi Zoltan)

• Allow `ecpg` to use noise words `FROM` and `IN` in `FETCH` and `MOVE` (Boszormenyi Zoltan)

E.158.3.10. Build Options

• Enable client thread safety by default (Bruce Momjian)

  The thread-safety option can be disabled with `configure --disable-thread-safety`.

• Add support for controlling the Linux out-of-memory killer (Alex Hunsaker, Tom Lane)

  Now that `/proc/self/oom_adj` allows disabling of the Linux out-of-memory (OOM) killer, it's recommendable to disable OOM kills for the postmaster. It may then be desirable to re-enable OOM kills for the postmaster's child processes. The new compile-time option `LINUX_OOM_ADJ` allows the killer to be reactivated for child processes.

E.158.3.10.1. Makefiles

• New Makefile targets `world`, `install-world`, and `installcheck-world` (Andrew Dunstan)

  These are similar to the existing `all`, `install`, and `installcheck` targets, but they also build the HTML documentation, build and test `contrib`, and test server-side languages and `ecpg`.

• Add data and documentation installation location control to PGXS Makefiles (Mark Cave-Ayland)

• Add Makefile rules to build the PostgreSQL documentation as a single HTML file or as a single plain-text file (Peter Eisentraut, Bruce Momjian)

E.158.3.10.2. Windows

• Support compiling on 64-bit Windows and running in 64-bit mode (Tsutomu Yamada, Magnus Hagander)

  This allows for large shared memory sizes on Windows.
• Support server builds using Visual Studio 2008 (Magnus Hagander)

E.158.3.11. Source Code

• Distribute prebuilt documentation in a subdirectory tree, rather than as tar archive files inside the distribution tarball (Peter Eisentraut)

  For example, the prebuilt HTML documentation is now in `doc/src/sgml/html/`; the manual pages are packaged similarly.

• Make the server’s lexer reentrant (Tom Lane)

  This was needed for use of the lexer by PL/pgSQL.

• Improve speed of memory allocation (Tom Lane, Greg Stark)

• User-defined constraint triggers now have entries in `pg_constraint` as well as `pg_trigger` (Tom Lane)

  Because of this change, `pg_constraint.pgconstrname` is now redundant and has been removed.

• Add system catalog columns `pg_constraint.conindid` and `pg_trigger.tgconstrid` to better document the use of indexes for constraint enforcement (Tom Lane)

• Allow multiple conditions to be communicated to backends using a single operating system signal (Fujii Masao)

  This allows new features to be added without a platform-specific constraint on the number of signal conditions.

• Improve source code test coverage, including `contrib`, PL/Python, and PL/Perl (Peter Eisentraut, Andrew Dunstan)

• Remove the use of flat files for system table bootstrapping (Tom Lane, Alvaro Herrera)

  This improves performance when using many roles or databases, and eliminates some possible failure conditions.

• Automatically generate the initial contents of `pg_attribute` for “bootstrapped” catalogs (John Naylor)

  This greatly simplifies changes to these catalogs.

• Split the processing of `INSERT/UPDATE/DELETE` operations out of `execMain.c` (Marko Tiikkaja)

  Updates are now executed in a separate ModifyTable node. This change is necessary infrastructure for future improvements.

• Simplify translation of `psql`’s SQL help text (Peter Eisentraut)

• Reduce the lengths of some file names so that all file paths in the distribution tarball are less than 100 characters (Tom Lane)

  Some decompression programs have problems with longer file paths.

• Add a new `ERRCODE_INVALID_PASSWORD` SQLSTATE error code (Bruce Momjian)

• With authors’ permissions, remove the few remaining personal source code copyright notices (Bruce Momjian)

2917
The personal copyright notices were insignificant but the community occasionally had to answer questions about them.

- Add new documentation section about running PostgreSQL in non-durable mode to improve performance (Bruce Momjian)
- Restructure the HTML documentation Makefile rules to make their dependency checks work correctly, avoiding unnecessary rebuilds (Peter Eisentraut)
- Use DocBook XSL stylesheets for man page building, rather than Docbook2X (Peter Eisentraut)
  This changes the set of tools needed to build the man pages.
- Improve PL/Perl code structure (Tim Bunce)
- Improve error context reports in PL/Perl (Alexey Klyukin)

E.158.3.11.1. New Build Requirements

Note that these requirements do not apply when building from a distribution tarball, since tarballs include the files that these programs are used to build.

- Require Autoconf 2.63 to build configure (Peter Eisentraut)
- Require Flex 2.5.31 or later to build from a CVS checkout (Tom Lane)
- Require Perl version 5.8 or later to build from a CVS checkout (John Naylor, Andrew Dunstan)

E.158.3.11.2. Portability

- Use a more modern API for Bonjour (Tom Lane)
  Bonjour support now requires macOS 10.3 or later. The older API has been deprecated by Apple.
- Add spinlock support for the SuperH architecture (Nobuhiro Iwamatsu)
- Allow non-GCC compilers to use inline functions if they support them (Kurt Harriman)
- Remove support for platforms that don't have a working 64-bit integer data type (Tom Lane)
- Restructure use of LDFLAGS to be more consistent across platforms (Tom Lane)
  LDFLAGS is now used for linking both executables and shared libraries, and we add on LD¬FLAGS_EX when linking executables, or LDFLAGS_SL when linking shared libraries.

E.158.3.11.3. Server Programming

- Make backend header files safe to include in C++ (Kurt Harriman, Peter Eisentraut)
  These changes remove keyword conflicts that previously made C++ usage difficult in backend code. However, there are still other complexities when using C++ for backend functions. extern "C" ( ) is still necessary in appropriate places, and memory management and error handling are still problematic.
- Add AggCheckCallContext() for use in detecting if a C function is being called as an aggregate (Hitoshi Harada)
- Change calling convention for SearchSysCache() and related functions to avoid hard-wiring the maximum number of cache keys (Robert Haas)

Existing calls will still work for the moment, but can be expected to break in 9.1 or later if not converted to the new style.
• Require calls of `fastgetattr()` and `heap_getattr()` backend macros to provide a non-NULL fourth argument (Robert Haas)

• Custom typanalyze functions should no longer rely on `VacAttrStats.attr` to determine the type of data they will be passed (Tom Lane)

    This was changed to allow collection of statistics on index columns for which the storage type is different from the underlying column data type. There are new fields that tell the actual datatype being analyzed.

### E.158.3.11.4. Server Hooks

• Add parser hooks for processing ColumnRef and ParamRef nodes (Tom Lane)

• Add a ProcessUtility hook so loadable modules can control utility commands (Itagaki Takahiro)

### E.158.3.11.5. Binary Upgrade Support

• Add `contrib/pg_upgrade` to support in-place upgrades (Bruce Momjian)

    This avoids the requirement of dumping/reloading the database when upgrading to a new major release of PostgreSQL, thus reducing downtime by orders of magnitude. It supports upgrades to 9.0 from PostgreSQL 8.3 and 8.4.

• Add support for preserving relation `relnode` values during binary upgrades (Bruce Momjian)

• Add support for preserving `pg_type` and `pg_enum` OIDs during binary upgrades (Bruce Momjian)

• Move data files within tablespaces into PostgreSQL-version-specific subdirectories (Bruce Momjian)

    This simplifies binary upgrades.

### E.158.3.12. Contrib

• Add multithreading option (`-j`) to `contrib/pgbench` (Itagaki Takahiro)

    This allows multiple CPUs to be used by pgbench, reducing the risk of pgbench itself becoming the test bottleneck.

• Add `\shell` and `\setshell` meta commands to `contrib/pgbench` (Michael Paquier)

• New features for `contrib/dict_xsyn` (Sergey Karpov)

    The new options are `matchorig`, `matchsynonyms`, and `keepsynonyms`.

• Add full text dictionary `contrib/unaccent` (Teodor Sigaev)

    This filtering dictionary removes accents from letters, which makes full-text searches over multiple languages much easier.

• Add `dblink_get_notify()` to `contrib/dblink` (Marcus Kempe)

    This allows asynchronous notifications in dblink.

• Improve `contrib/dblink`'s handling of dropped columns (Tom Lane)

    This affects `dblink_build_sql_insert()` and related functions. These functions now number columns according to logical not physical column numbers.
• Greatly increase contrib/hstore's data length limit, and add B-tree and hash support so GROUP BY and DISTINCT operations are possible on hstore columns (Andrew Gierth)

New functions and operators were also added. These improvements make hstore a full-function key-value store embedded in PostgreSQL.

• Add contrib/passwordcheck to support site-specific password strength policies (Laurenz Albe)

The source code of this module should be modified to implement site-specific password policies.

• Add contrib/pg_archivecleanup tool (Simon Riggs)

This is designed to be used in the archive_cleanup_command server parameter, to remove no-longer-needed archive files.

• Add query text to contrib/auto_explain output (Andrew Dunstan)

• Add buffer access counters to contrib/pg_stat_statements (Itagaki Takahiro)

• Update contrib/start-scripts/linux to use /proc/self/oom_adj to disable the Linux out-of-memory (OOM) killer (Alex Hunsaker, Tom Lane)

E.159. Release 8.4.22

Release date: 2014-07-24

This release contains a variety of fixes from 8.4.21. For information about new features in the 8.4 major release, see Section E.181.

This is expected to be the last PostgreSQL release in the 8.4.X series. Users are encouraged to update to a newer release branch soon.

E.159.1. Migration to Version 8.4.22

A dump/restore is not required for those running 8.4.X.

However, this release corrects an index corruption problem in some GiST indexes. See the first changelog entry below to find out whether your installation has been affected and what steps you should take if so.

Also, if you are upgrading from a version earlier than 8.4.19, see Section E.162.

E.159.2. Changes

• Correctly initialize padding bytes in contrib/btree_gist indexes on bit columns (Heikki Linnakangas)

This error could result in incorrect query results due to values that should compare equal not being seen as equal. Users with GiST indexes on bit or bit varying columns should REINDEX those indexes after installing this update.

• Protect against torn pages when deleting GIN list pages (Heikki Linnakangas)

This fix prevents possible index corruption if a system crash occurs while the page update is being written to disk.

• Fix possibly-incorrect cache invalidation during nested calls to ReceiveSharedInvalidMessages (Andres Freund)
• Don’t assume a subquery’s output is unique if there’s a set-returning function in its targetlist (David Rowley)

This oversight could lead to misoptimization of constructs like WHERE x IN (SELECT y, generate_series(1,10) FROM t GROUP BY y).

• Fix failure to detoast fields in composite elements of structured types (Tom Lane)

This corrects cases where TOAST pointers could be copied into other tables without being dereferenced. If the original data is later deleted, it would lead to errors like “missing chunk number 0 for toast value ...” when the now-dangling pointer is used.

• Fix “record type has not been registered” failures with whole-row references to the output of Append plan nodes (Tom Lane)

• Fix possible crash when invoking a user-defined function while rewinding a cursor (Tom Lane)

• Fix query-lifespan memory leak while evaluating the arguments for a function in FROM (Tom Lane)

• Fix session-lifespan memory leaks in regular-expression processing (Tom Lane, Arthur O’Dwyer, Greg Stark)

• Fix data encoding error in hungarian.stop (Tom Lane)

• Fix liveness checks for rows that were inserted in the current transaction and then deleted by a now-rolled-back subtransaction (Andres Freund)

This could cause problems (at least spurious warnings, and at worst an infinite loop) if CREATE INDEX or CLUSTER were done later in the same transaction.

• Clear pg_stat_activity.xact_start during PREPARE TRANSACTION (Andres Freund)

After the PREPARE, the originating session is no longer in a transaction, so it should not continue to display a transaction start time.

• Fix REASSIGN OWNED to not fail for text search objects (Álvaro Herrera)

• Block signals during postmaster startup (Tom Lane)

This ensures that the postmaster will properly clean up after itself if, for example, it receives SIGINT while still starting up.

• Secure Unix-domain sockets of temporary postmasters started during make check (Noah Misch)

Any local user able to access the socket file could connect as the server’s bootstrap superuser, then proceed to execute arbitrary code as the operating-system user running the test, as we previously noted in CVE-2014-0067. This change defends against that risk by placing the server’s socket in a temporary, mode 0700 subdirectory of /tmp. The hazard remains however on platforms where Unix sockets are not supported, notably Windows, because then the temporary postmaster must accept local TCP connections.

A useful side effect of this change is to simplify make check testing in builds that override DEFAULT_PGSOCKET_DIR. Popular non-default values like /var/run/postgresql are often not writable by the build user, requiring workarounds that will no longer be necessary.

• On Windows, allow new sessions to absorb values of PGC_BACKEND parameters (such as log_connections) from the configuration file (Amit Kapila)

Previously, if such a parameter were changed in the file post-startup, the change would have no effect.

• Properly quote executable path names on Windows (Nikhil Deshpande)
This oversight could cause initdb and pg_upgrade to fail on Windows, if the installation path contained both spaces and @ signs.

- Fix linking of libpython on macOS (Tom Lane)

  The method we previously used can fail with the Python library supplied by Xcode 5.0 and later.

- Avoid buffer bloat in libpq when the server consistently sends data faster than the client can absorb it (Shin-ichi Morita, Tom Lane)

  libpq could be coerced into enlarging its input buffer until it runs out of memory (which would be reported misleadingly as “lost synchronization with server”). Under ordinary circumstances it’s quite far-fetched that data could be continuously transmitted more quickly than the recv() loop can absorb it, but this has been observed when the client is artificially slowed by scheduler constraints.

- Ensure that LDAP lookup attempts in libpq time out as intended (Laurenz Albe)

- Fix pg_restore's processing of old-style large object comments (Tom Lane)

  A direct-to-database restore from an archive file generated by a pre-9.0 version of pg_dump would usually fail if the archive contained more than a few comments for large objects.

- In contrib/pgcrypto functions, ensure sensitive information is cleared from stack variables before returning (Marko Kreen)

- In contrib/uuid-ossp, cache the state of the OSSP UUID library across calls (Tom Lane)

  This improves the efficiency of UUID generation and reduces the amount of entropy drawn from /dev/urandom, on platforms that have that.

- Update time zone data files to tzdata release 2014e for DST law changes in Crimea, Egypt, and Morocco.

## E.160. Release 8.4.21

**Release date:** 2014-03-20

This release contains a variety of fixes from 8.4.20. For information about new features in the 8.4 major release, see Section E.181.

The PostgreSQL community will stop releasing updates for the 8.4.X release series in July 2014. Users are encouraged to update to a newer release branch soon.

### E.160.1. Migration to Version 8.4.21

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.19, see Section E.162.

### E.160.2. Changes

- Restore GIN metapages unconditionally to avoid torn-page risk (Heikki Linnakangas)

  Although this oversight could theoretically result in a corrupted index, it is unlikely to have caused any problems in practice, since the active part of a GIN metapage is smaller than a standard 512-byte disk sector.

- Allow regular-expression operators to be terminated early by query cancel requests (Tom Lane)
This prevents scenarios wherein a pathological regular expression could lock up a server process uninterruptibly for a long time.

• Remove incorrect code that tried to allow OVERLAPS with single-element row arguments (Joshua Yanovski)

This code never worked correctly, and since the case is neither specified by the SQL standard nor documented, it seemed better to remove it than fix it.

• Avoid getting more than AccessShareLock when de-parsing a rule or view (Dean Rasheed)

This oversight resulted in pg_dump unexpectedly acquiring RowExclusiveLock locks on tables mentioned as the targets of INSERT/UPDATE/DELETE commands in rules. While usually harmless, that could interfere with concurrent transactions that tried to acquire, for example, ShareLock on those tables.

• Prevent interrupts while reporting non-ERROR messages (Tom Lane)

This guards against rare server-process freezeups due to recursive entry to syslog(), and perhaps other related problems.

• Update time zone data files to tzdata release 2014a for DST law changes in Fiji and Turkey, plus historical changes in Israel and Ukraine.

**E.161. Release 8.4.20**

**Release date:** 2014-02-20

This release contains a variety of fixes from 8.4.19. For information about new features in the 8.4 major release, see Section E.181.

The PostgreSQL community will stop releasing updates for the 8.4.X release series in July 2014. Users are encouraged to update to a newer release branch soon.

**E.161.1. Migration to Version 8.4.20**

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.19, see Section E.162.

**E.161.2. Changes**

• Shore up GRANT ... WITH ADMIN OPTION restrictions (Noah Misch)

Granting a role without ADMIN OPTION is supposed to prevent the grantee from adding or removing members from the granted role, but this restriction was easily bypassed by doing SET ROLE first. The security impact is mostly that a role member can revoke the access of others, contrary to the wishes of his grantor. Unapproved role member additions are a lesser concern, since an uncooperative role member could provide most of his rights to others anyway by creating views or SECURITY DEFINER functions. (CVE-2014-0060)

• Prevent privilege escalation via manual calls to PL validator functions (Andres Freund)

The primary role of PL validator functions is to be called implicitly during CREATE FUNCTION, but they are also normal SQL functions that a user can call explicitly. Calling a validator on a function actually written in some other language was not checked for and could be exploited for privilege-escalation purposes. The fix involves adding a call to a privilege-checking function in each validator function. Non-core procedural languages will also need to make this change to their own validator functions, if any. (CVE-2014-0061)
• Avoid multiple name lookups during table and index DDL (Robert Haas, Andres Freund)

If the name lookups come to different conclusions due to concurrent activity, we might perform some parts of the DDL on a different table than other parts. At least in the case of `CREATE INDEX`, this can be used to cause the permissions checks to be performed against a different table than the index creation, allowing for a privilege escalation attack. (CVE-2014-0062)

• Prevent buffer overrun with long datetime strings (Noah Misch)

The `MAXDATELEN` constant was too small for the longest possible value of type `interval`, allowing a buffer overrun in `interval_out()`. Although the datetime input functions were more careful about avoiding buffer overrun, the limit was short enough to cause them to reject some valid inputs, such as input containing a very long timezone name. The ecpg library contained these vulnerabilities along with some of its own. (CVE-2014-0063)

• Prevent buffer overrun due to integer overflow in size calculations (Noah Misch, Heikki Linnakangas)

Several functions, mostly type input functions, calculated an allocation size without checking for overflow. If overflow did occur, a too-small buffer would be allocated and then written past. (CVE-2014-0064)

• Prevent overruns of fixed-size buffers (Peter Eisentraut, Jozef Mlich)

Use `strlcpy()` and related functions to provide a clear guarantee that fixed-size buffers are not overrun. Unlike the preceding items, it is unclear whether these cases really represent live issues, since in most cases there appear to be previous constraints on the size of the input string. Nonetheless it seems prudent to silence all Coverity warnings of this type. (CVE-2014-0065)

• Avoid crashing if `crypt()` returns NULL (Honza Horak, Bruce Momjian)

There are relatively few scenarios in which `crypt()` could return NULL, but `contrib/chkpass` would crash if it did. One practical case in which this could be an issue is if libc is configured to refuse to execute unapproved hashing algorithms (e.g., “FIPS mode”). (CVE-2014-0066)

• Document risks of `make check` in the regression testing instructions (Noah Misch, Tom Lane)

Since the temporary server started by `make check` uses “trust” authentication, another user on the same machine could connect to it as database superuser, and then potentially exploit the privileges of the operating-system user who started the tests. A future release will probably incorporate changes in the testing procedure to prevent this risk, but some public discussion is needed first. So for the moment, just warn people against using `make check` when there are untrusted users on the same machine. (CVE-2014-0067)

• Fix possible mis-replay of WAL records when some segments of a relation aren’t full size (Greg Stark, Tom Lane)

The WAL update could be applied to the wrong page, potentially many pages past where it should have been. Aside from corrupting data, this error has been observed to result in significant “bloat” of standby servers compared to their masters, due to updates being applied far beyond where the end-of-file should have been. This failure mode does not appear to be a significant risk during crash recovery, only when initially synchronizing a standby created from a base backup taken from a quickly-changing master.

• Ensure that insertions into non-leaf GIN index pages write a full-page WAL record when appropriate (Heikki Linnakangas)

The previous coding risked index corruption in the event of a partial-page write during a system crash.

• Fix race conditions during server process exit (Robert Haas)
Ensure that signal handlers don’t attempt to use the process’s MyProc pointer after it’s no longer valid.

- Fix unsafe references to errno within error reporting logic (Christian Kruse)

  This would typically lead to odd behaviors such as missing or inappropriate HINT fields.

- Fix possible crashes from using ereport() too early during server startup (Tom Lane)

  The principal case we’ve seen in the field is a crash if the server is started in a directory it doesn’t have permission to read.

- Clear retry flags properly in OpenSSL socket write function (Alexander Kukushkin)

  This omission could result in a server lockup after unexpected loss of an SSL-encrypted connection.

- Fix length checking for Unicode identifiers (U&"..." syntax) containing escapes (Tom Lane)

  A spurious truncation warning would be printed for such identifiers if the escaped form of the identifier was too long, but the identifier actually didn’t need truncation after de-escaping.

- Fix possible crash due to invalid plan for nested sub-selects, such as WHERE (... x IN (SELECT ...) ...) IN (SELECT ...) (Tom Lane)

- Ensure that ANALYZE creates statistics for a table column even when all the values in it are “too wide” (Tom Lane)

  ANALYZE intentionally omits very wide values from its histogram and most-common-values calculations, but it neglected to do something sane in the case that all the sampled entries are too wide.

- In ALTER TABLE ... SET TABLESPACE, allow the database's default tablespace to be used without a permissions check (Stephen Frost)

  CREATE TABLE has always allowed such usage, but ALTER TABLE didn't get the memo.

- Fix “cannot accept a set” error when some arms of a CASE return a set and others don’t (Tom Lane)

- Fix checks for all-zero client addresses in pgstat functions (Kevin Grittner)

- Fix possible misclassification of multibyte characters by the text search parser (Tom Lane)

  Non-ASCII characters could be misclassified when using C locale with a multibyte encoding. On Cygwin, non-C locales could fail as well.

- Fix possible misbehavior in plainto_tsquery() (Heikki Linnakangas)

  Use memmove() not memcpy() for copying overlapping memory regions. There have been no field reports of this actually causing trouble, but it’s certainly risky.

- Accept SHIFT_JIS as an encoding name for locale checking purposes (Tatsuo Ishii)

- Fix misbehavior of PQhost() on Windows (Fujii Masao)

  It should return localhost if no host has been specified.

- Improve error handling in libpq and psql for failures during COPY TO STDOUT/FROM STDIN (Tom Lane)

  In particular this fixes an infinite loop that could occur in 9.2 and up if the server connection was lost during COPY FROM STDIN. Variants of that scenario might be possible in older versions, or with other client applications.

- Fix misaligned descriptors in ecpg (MauMau)
• In ecpg, handle lack of a hostname in the connection parameters properly (Michael Meskes)

• Fix performance regression in contrib/dblink connection startup (Joe Conway)
  Avoid an unnecessary round trip when client and server encodings match.

• In contrib/ism, fix incorrect calculation of the check digit for ISMN values (Fabien Coelho)

• Ensure client-code-only installation procedure works as documented (Peter Eisentraut)

• In Mingw and Cygwin builds, install the libpq DLL in the bin directory (Andrew Dunstan)
  This duplicates what the MSVC build has long done. It should fix problems with programs like psql failing to start because they can't find the DLL.

• Don't generate plain-text HISTORY and src/test/regress/README files anymore (Tom Lane)
  These text files duplicated the main HTML and PDF documentation formats. The trouble involved in maintaining them greatly outweighs the likely audience for plain-text format. Distribution tarballs will still contain files by these names, but they'll just be stubs directing the reader to consult the main documentation. The plain-text INSTALL file will still be maintained, as there is arguably a use-case for that.

• Update time zone data files to tzdata release 2013i for DST law changes in Jordan and historical changes in Cuba.
  In addition, the zones Asia/Riyadh87, Asia/Riyadh88, and Asia/Riyadh89 have been removed, as they are no longer maintained by IANA, and never represented actual civil timekeeping practice.

E.162. Release 8.4.19

Release date: 2013-12-05

This release contains a variety of fixes from 8.4.18. For information about new features in the 8.4 major release, see Section E.181.

E.162.1. Migration to Version 8.4.19

A dump/restore is not required for those running 8.4.X.

However, this release corrects a potential data corruption issue. See the first changelog entry below to find out whether your installation has been affected and what steps you can take if so.

Also, if you are upgrading from a version earlier than 8.4.17, see Section E.164.

E.162.2. Changes

• Fix VACUUM's tests to see whether it can update relfrozenxid (Andres Freund)
  In some cases VACUUM (either manual or autovacuum) could incorrectly advance a table's relfrozenxid value, allowing tuples to escape freezing, causing those rows to become invisible once $2^{31}$ transactions have elapsed. The probability of data loss is fairly low since multiple incorrect advancements would need to happen before actual loss occurs, but it's not zero. Users upgrading from release 8.4.8 or earlier are not affected, but all later versions contain the bug.

  The issue can be ameliorated by, after upgrading, vacuuming all tables in all databases while having vacuum_freeze_table_age set to zero. This will fix any latent corruption but will not be able to fix all pre-existing data errors. However, an installation can be presumed safe after performing
this vacuuming if it has executed fewer than 2^31 update transactions in its lifetime (check this with SELECT txid_current() < 2^31).

- Fix race condition in GIN index posting tree page deletion (Heikki Linnakangas)

  This could lead to transient wrong answers or query failures.

- Avoid flattening a subquery whose SELECT list contains a volatile function wrapped inside a sub-SELECT (Tom Lane)

  This avoids unexpected results due to extra evaluations of the volatile function.

- Fix planner's processing of non-simple-variable subquery outputs nested within outer joins (Tom Lane)

  This error could lead to incorrect plans for queries involving multiple levels of subqueries within JOIN syntax.

- Fix premature deletion of temporary files (Andres Freund)

- Fix possible read past end of memory in rule printing (Peter Eisentraut)

- Fix array slicing of int2vector and oidvector values (Tom Lane)

  Expressions of this kind are now implicitly promoted to regular int2 or oid arrays.

- Fix incorrect behaviors when using a SQL-standard, simple GMT offset timezone (Tom Lane)

  In some cases, the system would use the simple GMT offset value when it should have used the regular timezone setting that had prevailed before the simple offset was selected. This change also causes the timeofday function to honor the simple GMT offset zone.

- Prevent possible misbehavior when logging translations of Windows error codes (Tom Lane)

- Properly quote generated command lines in pg_ctl (Naoya Anzai and Tom Lane)

  This fix applies only to Windows.

- Fix pg_dumpall to work when a source database sets default_transaction_read_only via ALTER DATABASE SET (Kevin Grittner)

  Previously, the generated script would fail during restore.

- Fix ecpg's processing of lists of variables declared varchar (Zoltán Böszörményi)

- Make contrib/lo defend against incorrect trigger definitions (Marc Cousin)

- Update time zone data files to tzdata release 2013h for DST law changes in Argentina, Brazil, Jordan, Libya, Liechtenstein, Morocco, and Palestine. Also, new timezone abbreviations WIB, WIT, WITA for Indonesia.

## E.163. Release 8.4.18

**Release date:** 2013-10-10

This release contains a variety of fixes from 8.4.17. For information about new features in the 8.4 major release, see Section E.181.

### E.163.1. Migration to Version 8.4.18

A dump/restore is not required for those running 8.4.X.
However, if you are upgrading from a version earlier than 8.4.17, see Section E.164.

E.163.2. Changes

• Prevent corruption of multi-byte characters when attempting to case-fold identifiers (Andrew Dunstan)

  PostgreSQL case-folds non-ASCII characters only when using a single-byte server encoding.

• Fix memory leak caused by lo_open() failure (Heikki Linnakangas)

• Fix memory overcommit bug when work_mem is using more than 24GB of memory (Stephen Frost)

• Fix deadlock bug in libpq when using SSL (Stephen Frost)

• Properly compute row estimates for boolean columns containing many NULL values (Andrew Gierth)

  Previously tests like col IS NOT TRUE and col IS NOT FALSE did not properly factor in NULL values when estimating plan costs.

• Prevent pushing down WHERE clauses into unsafe UNION/INTERSECT subqueries (Tom Lane)

  Subqueries of a UNION or INTERSECT that contain set-returning functions or volatile functions in their SELECT lists could be improperly optimized, leading to run-time errors or incorrect query results.

• Fix rare case of “failed to locate grouping columns” planner failure (Tom Lane)

• Improve view dumping code’s handling of dropped columns in referenced tables (Tom Lane)

• Fix possible deadlock during concurrent CREATE INDEX CONCURRENTLY operations (Tom Lane)

• Fix regexp_matches() handling of zero-length matches (Jeevan Chalke)

  Previously, zero-length matches like ‘~’ could return too many matches.

• Fix crash for overly-complex regular expressions (Heikki Linnakangas)

• Fix regular expression match failures for back references combined with non-greedy quantifiers (Jeevan Chalke)

• Prevent CREATE FUNCTION from checking SET variables unless function body checking is enabled (Tom Lane)

• Fix pgp_pub_decrypt() so it works for secret keys with passwords (Marko Kreen)

• Remove rare inaccurate warning during vacuum of index-less tables (Heikki Linnakangas)

• Avoid possible failure when performing transaction control commands (e.g ROLLBACK) in prepared queries (Tom Lane)

• Ensure that floating-point data input accepts standard spellings of “infinity” on all platforms (Tom Lane)

  The C99 standard says that allowable spellings are inf, +inf, -inf, infinity, +infinity, and -infinity. Make sure we recognize these even if the platform’s strtod function doesn’t.

• Expand ability to compare rows to records and arrays (Rafal Rzepecki, Tom Lane)
• Update time zone data files to tzdata release 2013d for DST law changes in Israel, Morocco, Palestine, and Paraguay. Also, historical zone data corrections for Macquarie Island.

E.164. Release 8.4.17

Release date: 2013-04-04

This release contains a variety of fixes from 8.4.16. For information about new features in the 8.4 major release, see Section E.181.

E.164.1. Migration to Version 8.4.17

A dump/restore is not required for those running 8.4.X. However, this release corrects several errors in management of GiST indexes. After installing this update, it is advisable to REINDEX any GiST indexes that meet one or more of the conditions described below.

Also, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

E.164.2. Changes

• Reset OpenSSL randomness state in each postmaster child process (Marko Kreen)

This avoids a scenario wherein random numbers generated by contrib/pgcrypto functions might be relatively easy for another database user to guess. The risk is only significant when the postmaster is configured with ssl = on but most connections don't use SSL encryption. (CVE-2013-1900)

• Fix GiST indexes to not use “fuzzy” geometric comparisons when it’s not appropriate to do so (Alexander Korotkov)

The core geometric types perform comparisons using “fuzzy” equality, but gist_box_same must do exact comparisons, else GiST indexes using it might become inconsistent. After installing this update, users should REINDEX any GiST indexes on box, polygon, circle, or point columns, since all of these use gist_box_same.

• Fix erroneous range-union and penalty logic in GiST indexes that use contrib/btree_gist for variable-width data types, that is text, bytea, bit, and numeric columns (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in useless index bloat. Users are advised to REINDEX such indexes after installing this update.

• Fix bugs in GiST page splitting code for multi-column indexes (Tom Lane)

These errors could result in inconsistent indexes in which some keys that are present would not be found by searches, and also in indexes that are unnecessarily inefficient to search. Users are advised to REINDEX multi-column GiST indexes after installing this update.

• Fix infinite-loop risk in regular expression compilation (Tom Lane, Don Porter)

• Fix potential null-pointer dereference in regular expression compilation (Tom Lane)

• Fix to_char() to use ASCII-only case-folding rules where appropriate (Tom Lane)

This fixes misbehavior of some template patterns that should be locale-independent, but mishandled “I” and “i” in Turkish locales.

• Fix unwanted rejection of timestamp 1999-12-31 24:00:00 (Tom Lane)
• Remove useless “picksplit doesn’t support secondary split” log messages (Josh Hansen, Tom Lane)

This message seems to have been added in expectation of code that was never written, and probably never will be, since GiST’s default handling of secondary splits is actually pretty good. So stop nagging end users about it.

• Fix possible failure to send a session’s last few transaction commit/abort counts to the statistics collector (Tom Lane)

• Eliminate memory leaks in PL/Perl's `spi_prepare()` function (Alex Hunsaker, Tom Lane)

• Fix `pg_dumpall` to handle database names containing “=” correctly (Heikki Linnakangas)

• Avoid crash in `pg_dump` when an incorrect connection string is given (Heikki Linnakangas)

• Ignore invalid indexes in `pg_dump` (Michael Paquier)

Dumping invalid indexes can cause problems at restore time, for example if the reason the index creation failed was because it tried to enforce a uniqueness condition not satisfied by the table's data. Also, if the index creation is in fact still in progress, it seems reasonable to consider it to be an uncommitted DDL change, which `pg_dump` wouldn't be expected to dump anyway.

• Fix `contrib/pg_trgm`'s `similarity()` function to return zero for trigram-less strings (Tom Lane)

Previously it returned `NaN` due to internal division by zero.

• Update time zone data files to `tzdata` release 2013b for DST law changes in Chile, Haiti, Morocco, Paraguay, and some Russian areas. Also, historical zone data corrections for numerous places.

Also, update the time zone abbreviation files for recent changes in Russia and elsewhere: CHOT, GET, IRKT, KGT, KRAT, MAGT, MAWT, MSK, NOVT, OGMT, TKT, VLAT, WST, YAKT, YEKT now follow their current meanings, and VOLT (Europe/Volgograd) and MIST (Antarctica/Macquarie) are added to the default abbreviations list.

E.165. Release 8.4.16

Release date: 2013-02-07

This release contains a variety of fixes from 8.4.15. For information about new features in the 8.4 major release, see Section E.181.

E.165.1. Migration to Version 8.4.16

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

E.165.2. Changes

• Prevent execution of `enum_recv` from SQL (Tom Lane)

The function was misdeclared, allowing a simple SQL command to crash the server. In principle an attacker might be able to use it to examine the contents of server memory. Our thanks to Sumit Soni (via Secunia SVCRP) for reporting this issue. (CVE-2013-0255)

• Update minimum recovery point when truncating a relation file (Heikki Linnakangas)

Once data has been discarded, it's no longer safe to stop recovery at an earlier point in the timeline.
• Fix SQL grammar to allow subscripting or field selection from a sub-SELECT result (Tom Lane)

• Protect against race conditions when scanning pg_tablespace (Stephen Frost, Tom Lane)

CREATE DATABASE and DROP DATABASE could misbehave if there were concurrent updates of pg_tablespace entries.

• Prevent DROP OWNED from trying to drop whole databases or tablespaces (Álvaro Herrera)

For safety, ownership of these objects must be reassigned, not dropped.

• Fix error in vacuum_freeze_table_age implementation (Andres Freund)

In installations that have existed for more than vacuum_freeze_min_age transactions, this mistake prevented autovacuum from using partial-table scans, so that a full-table scan would always happen instead.

• Prevent misbehavior when a RowExpr or XmlExpr is parse-analyzed twice (Andres Freund, Tom Lane)

This mistake could be user-visible in contexts such as CREATE TABLE LIKE INCLUDING INDEXES.

• Improve defenses against integer overflow in hashtable sizing calculations (Jeff Davis)

• Reject out-of-range dates in to_date() (Hitoshi Harada)

• Ensure that non-ASCII prompt strings are translated to the correct code page on Windows (Alexander Law, Noah Misch)

This bug affected psql and some other client programs.

• Fix possible crash in psql's \? command when not connected to a database (Meng Qingzhong)

• Fix one-byte buffer overrun in libpq’s PQprintTuples (Xi Wang)

This ancient function is not used anywhere by PostgreSQL itself, but it might still be used by some client code.

• Make ecpglib use translated messages properly (Chen Huajun)

• Properly install ecpg_comapt and pgtypes libraries on MSVC (Jiang Guiqing)

• Rearrange configure's tests for supplied functions so it is not fooled by bogus exports from libedit/libreadline (Christoph Berg)

• Ensure Windows build number increases over time (Magnus Hagander)

• Make pgxs build executables with the right .exe suffix when cross-compiling for Windows (Zoltan Boszormenyi)

• Add new timezone abbreviation FET (Tom Lane)

This is now used in some eastern-European time zones.

E.166. Release 8.4.15

Release date: 2012-12-06

This release contains a variety of fixes from 8.4.14. For information about new features in the 8.4 major release, see Section E.181.
E.166.1. Migration to Version 8.4.15

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

E.166.2. Changes

- Fix multiple bugs associated with CREATE INDEX CONCURRENTLY (Andres Freund, Tom Lane)

  Fix CREATE INDEX CONCURRENTLY to use in-place updates when changing the state of an index's pg_index row. This prevents race conditions that could cause concurrent sessions to miss updating the target index, thus resulting in corrupt concurrently-created indexes.

  Also, fix various other operations to ensure that they ignore invalid indexes resulting from a failed CREATE INDEX CONCURRENTLY command. The most important of these is VACUUM, because an auto-vacuum could easily be launched on the table before corrective action can be taken to fix or remove the invalid index.

- Avoid corruption of internal hash tables when out of memory (Hitoshi Harada)

- Fix planning of non-strict equivalence clauses above outer joins (Tom Lane)

  The planner could derive incorrect constraints from a clause equating a non-strict construct to something else, for example WHERE COALESCE(foo, 0) = 0 when foo is coming from the nullable side of an outer join.

- Improve planner's ability to prove exclusion constraints from equivalence classes (Tom Lane)

- Fix partial-row matching in hashed subplans to handle cross-type cases correctly (Tom Lane)

  This affects multicolumn NOT IN subplans, such as WHERE (a, b) NOT IN (SELECT x, y FROM ...) when for instance b and y are int4 and int8 respectively. This mistake led to wrong answers or crashes depending on the specific datatypes involved.

- Acquire buffer lock when re-fetching the old tuple for an AFTER ROW UPDATE/DELETE trigger (Andres Freund)

  In very unusual circumstances, this oversight could result in passing incorrect data to the precheck logic for a foreign-key enforcement trigger. That could result in a crash, or in an incorrect decision about whether to fire the trigger.

- Fix ALTER COLUMN TYPE to handle inherited check constraints properly (Pavan Deolasee)

  This worked correctly in pre-8.4 releases, and now works correctly in 8.4 and later.

- Fix REASSIGN OWNED to handle grants on table spaces (Álvaro Herrera)

- Ignore incorrect pg_attribute entries for system columns for views (Tom Lane)

  Views do not have any system columns. However, we forgot to remove such entries when converting a table to a view. That's fixed properly for 9.3 and later, but in previous branches we need to defend against existing mis-converted views.

- Fix rule printing to dump INSERT INTO table DEFAULT VALUES correctly (Tom Lane)

- Guard against stack overflow when there are too many UNION/INTERSECT/EXCEPT clauses in a query (Tom Lane)

- Prevent platform-dependent failures when dividing the minimum possible integer value by -1 (Xi Wang, Tom Lane)
• Fix possible access past end of string in date parsing (Hitoshi Harada)

• Produce an understandable error message if the length of the path name for a Unix-domain socket exceeds the platform-specific limit (Tom Lane, Andrew Dunstan)

  Formerly, this would result in something quite unhelpful, such as "Non-recoverable failure in name resolution".

• Fix memory leaks when sending composite column values to the client (Tom Lane)

• Make pg_ctl more robust about reading the postmaster.pid file (Heikki Linnakangas)

  Fix race conditions and possible file descriptor leakage.

• Fix possible crash in psql if incorrectly-encoded data is presented and the client_encoding setting is a client-only encoding, such as SJIS (Jiang Guiqing)

• Fix bugs in the restore.sql script emitted by pg_dump in tar output format (Tom Lane)

  The script would fail outright on tables whose names include upper-case characters. Also, make the script capable of restoring data in --inserts mode as well as the regular COPY mode.

• Fix pg_restore to accept POSIX-conformant tar files (Brian Weaver, Tom Lane)

  The original coding of pg_dump's tar output mode produced files that are not fully conformant with the POSIX standard. This has been corrected for version 9.3. This patch updates previous branches so that they will accept both the incorrect and the corrected formats, in hopes of avoiding compatibility problems when 9.3 comes out.

• Fix pg_resetxlog to locate postmaster.pid correctly when given a relative path to the data directory (Tom Lane)

  This mistake could lead to pg_resetxlog not noticing that there is an active postmaster using the data directory.

• Fix libpq's lo_import() and lo_export() functions to report file I/O errors properly (Tom Lane)

• Fix ecpg's processing of nested structure pointer variables (Muhammad Usama)

• Make contrib/pageinspect's btree page inspection functions take buffer locks while examining pages (Tom Lane)

• Fix pgxs support for building loadable modules on AIX (Tom Lane)

  Building modules outside the original source tree didn't work on AIX.

• Update time zone data files to tzdata release 2012j for DST law changes in Cuba, Israel, Jordan, Libya, Palestine, Western Samoa, and portions of Brazil.

---

E.167. Release 8.4.14

Release date: 2012-09-24

This release contains a variety of fixes from 8.4.13. For information about new features in the 8.4 major release, see Section E.181.

E.167.1. Migration to Version 8.4.14

A dump/restore is not required for those running 8.4.X.
However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

**E.167.2. Changes**

- Fix planner's assignment of executor parameters, and fix executor's rescan logic for CTE plan nodes (Tom Lane)
  
  These errors could result in wrong answers from queries that scan the same WITH subquery multiple times.

- Improve page-splitting decisions in GiST indexes (Alexander Korotkov, Robert Haas, Tom Lane)
  
  Multi-column GiST indexes might suffer unexpected bloat due to this error.

- Fix cascading privilege revoke to stop if privileges are still held (Tom Lane)
  
  If we revoke a grant option from some role $X$, but $X$ still holds that option via a grant from someone else, we should not recursively revoke the corresponding privilege from role(s) $Y$ that $X$ had granted it to.

- Fix handling of SIGFPE when PL/Perl is in use (Andres Freund)
  
  Perl resets the process's SIGFPE handler to SIG_IGN, which could result in crashes later on. Restore the normal Postgres signal handler after initializing PL/Perl.

- Prevent PL/Perl from crashing if a recursive PL/Perl function is redefined while being executed (Tom Lane)

- Work around possible misoptimization in PL/Perl (Tom Lane)
  
  Some Linux distributions contain an incorrect version of pthread.h that results in incorrect compiled code in PL/Perl, leading to crashes if a PL/Perl function calls another one that throws an error.

- Update time zone data files to tzdata release 2012f for DST law changes in Fiji

**E.168. Release 8.4.13**

**Release date:** 2012-08-17

This release contains a variety of fixes from 8.4.12. For information about new features in the 8.4 major release, see Section E.181.

**E.168.1. Migration to Version 8.4.13**

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

**E.168.2. Changes**

- Prevent access to external files/URLs via XML entity references (Noah Misch, Tom Lane)
  
  `xml_parse()` would attempt to fetch external files or URLs as needed to resolve DTD and entity references in an XML value, thus allowing unprivileged database users to attempt to fetch data with the privileges of the database server. While the external data wouldn't get returned directly to the user, portions of it could be exposed in error messages if the data didn't parse as valid XML; and in any case the mere ability to check existence of a file might be useful to an attacker. (CVE-2012-3489)
• Prevent access to external files/URLs via contrib/xml2's xslt_process() (Peter Eisen- traut)

libxslt offers the ability to read and write both files and URLs through stylesheet commands, thus allowing unprivileged database users to both read and write data with the privileges of the database server. Disable that through proper use of libxslt's security options. (CVE-2012-3488)

Also, remove xslt_process()'s ability to fetch documents and stylesheets from external files/ URLs. While this was a documented “feature”, it was long regarded as a bad idea. The fix for CVE-2012-3489 broke that capability, and rather than expend effort on trying to fix it, we're just going to summarily remove it.

• Prevent too-early recycling of btree index pages (Noah Misch)

When we allowed read-only transactions to skip assigning XIDs, we introduced the possibility that a deleted btree page could be recycled while a read-only transaction was still in flight to it. This would result in incorrect index search results. The probability of such an error occurring in the field seems very low because of the timing requirements, but nonetheless it should be fixed.

• Fix crash-safety bug with newly-created-or-reset sequences (Tom Lane)

If ALTER SEQUENCE was executed on a freshly created or reset sequence, and then precisely one nextval() call was made on it, and then the server crashed, WAL replay would restore the sequence to a state in which it appeared that no nextval() had been done, thus allowing the first sequence value to be returned again by the next nextval() call. In particular this could manifest for serial columns, since creation of a serial column's sequence includes an ALTER SEQUENCE OWNED BY step.

• Ensure the backup_label file is fsync'd after pg_start_backup() (Dave Kerr)

• Back-patch 9.1 improvement to compress the fsync request queue (Robert Haas)

This improves performance during checkpoints. The 9.1 change has now seen enough field testing to seem safe to back-patch.

• Only allow autovacuum to be auto-canceled by a directly blocked process (Tom Lane)

The original coding could allow inconsistent behavior in some cases; in particular, an autovacuum could get canceled after less than deadlock_timeout grace period.

• Improve logging of autovacuum cancels (Robert Haas)

• Fix log collector so that log_truncate_on_rotation works during the very first log rotation after server start (Tom Lane)

• Fix WITH attached to a nested set operation (UNION/INTERSECT/EXCEPT) (Tom Lane)

• Ensure that a whole-row reference to a subquery doesn't include any extra GROUP BY or ORDER BY columns (Tom Lane)

• Disallow copying whole-row references in CHECK constraints and index definitions during CREATE TABLE (Tom Lane)

This situation can arise in CREATE TABLE with LIKE or INHERITS. The copied whole-row variable was incorrectly labeled with the row type of the original table not the new one. Rejecting the case seems reasonable for LIKE, since the row types might well diverge later. For INHERITS we should ideally allow it, with an implicit coercion to the parent table's row type; but that will require more work than seems safe to back-patch.

• Fix memory leak in ARRAY(SELECT ...) subqueries (Heikki Linnakangas, Tom Lane)

• Fix extraction of common prefixes from regular expressions (Tom Lane)
The code could get confused by quantified parenthesized subexpressions, such as \(^{(foo)}\)?bar. This would lead to incorrect index optimization of searches for such patterns.

- Fix bugs with parsing signed \(hh:mm\) and \(hh:mm:ss\) fields in interval constants (Amit Kapila, Tom Lane)
- Report errors properly in contrib/xml2's \texttt{xslt_process()} (Tom Lane)
- Update time zone data files to tzdata release 2012e for DST law changes in Morocco and Tokelau

**E.169. Release 8.4.12**

**Release date:** 2012-06-04

This release contains a variety of fixes from 8.4.11. For information about new features in the 8.4 major release, see Section E.181.

**E.169.1. Migration to Version 8.4.12**

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

**E.169.2. Changes**

- Fix incorrect password transformation in contrib/pgcrypto's \texttt{DES crypt()} function (Solar Designer)
  
  If a password string contained the byte value \(0x80\), the remainder of the password was ignored, causing the password to be much weaker than it appeared. With this fix, the rest of the string is properly included in the DES hash. Any stored password values that are affected by this bug will thus no longer match, so the stored values may need to be updated. (CVE-2012-2143)

- Ignore \texttt{SECURITY DEFINER} and \texttt{SET} attributes for a procedural language's call handler (Tom Lane)
  
  Applying such attributes to a call handler could crash the server. (CVE-2012-2655)

- Allow numeric timezone offsets in \texttt{timestamp} input to be up to 16 hours away from UTC (Tom Lane)
  
  Some historical time zones have offsets larger than 15 hours, the previous limit. This could result in dumped data values being rejected during reload.

- Fix timestamp conversion to cope when the given time is exactly the last DST transition time for the current timezone (Tom Lane)
  
  This oversight has been there a long time, but was not noticed previously because most DST-using zones are presumed to have an indefinite sequence of future DST transitions.

- Fix \texttt{text to name} and \texttt{char to name} casts to perform string truncation correctly in multibyte encodings (Karl Schnaitter)

- Fix memory copying bug in \texttt{to_tsquery()} (Heikki Linnakangas)

- Fix planner's handling of outer PlaceHolderVars within subqueries (Tom Lane)

  This bug concerns sub-SELECTs that reference variables coming from the nullable side of an outer join of the surrounding query. In 9.1, queries affected by this bug would fail with “ERROR: Up-
per-level PlaceHolderVar found where not expected". But in 9.0 and 8.4, you'd silently get possibly-wrong answers, since the value transmitted into the subquery wouldn't go to null when it should.

- Fix slow session startup when pg_attribute is very large (Tom Lane)

If pg_attribute exceeds one-fourth of shared_buffers, cache rebuilding code that is sometimes needed during session start would trigger the synchronized-scan logic, causing it to take many times longer than normal. The problem was particularly acute if many new sessions were starting at once.

- Ensure sequential scans check for query cancel reasonably often (Merlin Moncure)

A scan encountering many consecutive pages that contain no live tuples would not respond to interrupts meanwhile.

- Ensure the Windows implementation of PGSemaphoreLock() clears ImmediateInterruptOK before returning (Tom Lane)

This oversight meant that a query-cancel interrupt received later in the same query could be accepted at an unsafe time, with unpredictable but not good consequences.

- Ensure whole-row variables safely when printing views or rules (Abbas Butt, Tom Lane)

Corner cases involving ambiguous names (that is, the name could be either a table or column name of the query) were printed in an ambiguous way, risking that the view or rule would be interpreted differently after dump and reload. Avoid the ambiguous case by attaching a no-op cast.

- Fix COPY FROM to properly handle null marker strings that correspond to invalid encoding (Tom Lane)

A null marker string such as E\0 should work, and did work in the past, but the case got broken in 8.4.

- Ensure autovacuum worker processes perform stack depth checking properly (Heikki Linnakangas)

Previously, infinite recursion in a function invoked by auto-ANALYZE could crash worker processes.

- Fix logging collector to not lose log coherency under high load (Andrew Dunstan)

The collector previously could fail to reassemble large messages if it got too busy.

- Fix logging collector to ensure it will restart file rotation after receiving SIGHUP (Tom Lane)

- Fix WAL replay logic for GIN indexes to not fail if the index was subsequently dropped (Tom Lane)

- Fix memory leak in PL/pgSQL's RETURN NEXT command (Joe Conway)

- Fix PL/pgSQL's GET DIAGNOSTICS command when the target is the function's first variable (Tom Lane)

- Fix potential access off the end of memory in psql's expanded display (\x) mode (Peter Eisentraut)

- Fix several performance problems in pg_dump when the database contains many objects (Jeff Janes, Tom Lane)

pg_dump could get very slow if the database contained many schemas, or if many objects are in dependency loops, or if there are many owned sequences.

- Fix contrib/dblink's dblink_exec() to not leak temporary database connections upon error (Tom Lane)
• Fix contrib/dblink to report the correct connection name in error messages (Kyotaro Horiguchi)

• Update time zone data files to tzdata release 2012c for DST law changes in Antarctica, Armenia, Chile, Cuba, Falkland Islands, Gaza, Haiti, Hebron, Morocco, Syria, and Tokelau Islands; also historical corrections for Canada.

E.170. Release 8.4.11

Release date: 2012-02-27

This release contains a variety of fixes from 8.4.10. For information about new features in the 8.4 major release, see Section E.181.

E.170.1. Migration to Version 8.4.11

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.10, see Section E.171.

E.170.2. Changes

• Require execute permission on the trigger function for CREATE TRIGGER (Robert Haas)

  This missing check could allow another user to execute a trigger function with forged input data, by installing it on a table he owns. This is only of significance for trigger functions marked SECURITY DEFINER, since otherwise trigger functions run as the table owner anyway. (CVE-2012-0866)

• Remove arbitrary limitation on length of common name in SSL certificates (Heikki Linnakangas)

  Both libpq and the server truncated the common name extracted from an SSL certificate at 32 bytes. Normally this would cause nothing worse than an unexpected verification failure, but there are some rather-implausible scenarios in which it might allow one certificate holder to impersonate another. The victim would have to have a common name exactly 32 bytes long, and the attacker would have to persuade a trusted CA to issue a certificate in which the common name has that string as a prefix. Impersonating a server would also require some additional exploit to redirect client connections. (CVE-2012-0867)

• Convert newlines to spaces in names written in pg_dump comments (Robert Haas)

  pg_dump was incautious about sanitizing object names that are emitted within SQL comments in its output script. A name containing a newline would at least render the script syntactically incorrect. Maliciously crafted object names could present a SQL injection risk when the script is reloaded. (CVE-2012-0868)

• Fix btree index corruption from insertions concurrent with vacuuming (Tom Lane)

  An index page split caused by an insertion could sometimes cause a concurrently-running VACUUM to miss removing index entries that it should remove. After the corresponding table rows are removed, the dangling index entries would cause errors (such as "could not read block N in file ...") or worse, silently wrong query results after unrelated rows are re-inserted at the now-free table locations. This bug has been present since release 8.2, but occurs so infrequently that it was not diagnosed until now. If you have reason to suspect that it has happened in your database, reindexing the affected index will fix things.

• Update per-column permissions, not only per-table permissions, when changing table owner (Tom Lane)
Failure to do this meant that any previously granted column permissions were still shown as having been granted by the old owner. This meant that neither the new owner nor a superuser could revoke the now-untraceable-to-table-owner permissions.

- Allow non-existent values for some settings in ALTER USER/DATABASE SET (Heikki Linnakangas)

  Allow default_text_search_config, default_tablespace, and temp_tablespaces to be set to names that are not known. This is because they might be known in another database where the setting is intended to be used, or for the tablespace cases because the tablespace might not be created yet. The same issue was previously recognized for search_path, and these settings now act like that one.

- Avoid crashing when we have problems deleting table files post-commit (Tom Lane)

  Dropping a table should lead to deleting the underlying disk files only after the transaction commits. In event of failure then (for instance, because of wrong file permissions) the code is supposed to just emit a warning message and go on, since it’s too late to abort the transaction. This logic got broken as of release 8.4, causing such situations to result in a PANIC and an unrestartable database.

- Track the OID counter correctly during WAL replay, even when it wraps around (Tom Lane)

  Previously the OID counter would remain stuck at a high value until the system exited replay mode. The practical consequences of that are usually nil, but there are scenarios wherein a standby server that's been promoted to master might take a long time to advance the OID counter to a reasonable value once values are needed.

- Fix regular expression back-references with * attached (Tom Lane)

  Rather than enforcing an exact string match, the code would effectively accept any string that satisfies the pattern sub-expression referenced by the back-reference symbol.

  A similar problem still afflicts back-references that are embedded in a larger quantified expression, rather than being the immediate subject of the quantifier. This will be addressed in a future PostgreSQL release.

- Fix recently-introduced memory leak in processing of inet/cidr values (Heikki Linnakangas)

  A patch in the December 2011 releases of PostgreSQL caused memory leakage in these operations, which could be significant in scenarios such as building a btree index on such a column.

- Fix dangling pointer after CREATE TABLE AS/SELECT INTO in a SQL-language function (Tom Lane)

  In most cases this only led to an assertion failure in assert-enabled builds, but worse consequences seem possible.

- Avoid double close of file handle in syslogger on Windows (MauMau)

  Ordinarily this error was invisible, but it would cause an exception when running on a debug version of Windows.

- Fix I/O-conversion-related memory leaks in plpgsql (Andres Freund, Jan Urbanski, Tom Lane)

  Certain operations would leak memory until the end of the current function.

- Improve pg_dump's handling of inherited table columns (Tom Lane)

  pg_dump mishandled situations where a child column has a different default expression than its parent column. If the default is textually identical to the parent's default, but not actually the same
(for instance, because of schema search path differences) it would not be recognized as different, so that after dump and restore the child would be allowed to inherit the parent's default. Child columns that are NOT NULL where their parent is not could also be restored subtly incorrectly.

- Fix pg_restore's direct-to-database mode for INSERT-style table data (Tom Lane)

Direct-to-database restores from archive files made with --inserts or --column-inserts options fail when using pg_restore from a release dated September or December 2011, as a result of an oversight in a fix for another problem. The archive file itself is not at fault, and text-mode output is okay.

- Allow AT option in ecpg DEALLOCATE statements (Michael Meskes)

The infrastructure to support this has been there for awhile, but through an oversight there was still an error check rejecting the case.

- Fix error in contrib/intarray's int[] & int[] operator (Guillaume Lelarge)

If the smallest integer the two input arrays have in common is 1, and there are smaller values in either array, then 1 would be incorrectly omitted from the result.

- Fix error detection in contrib/pgcrypto's encrypt_iv() and decrypt_iv() (Marko Kreen)

These functions failed to report certain types of invalid-input errors, and would instead return random garbage values for incorrect input.

- Fix one-byte buffer overrun in contrib/test_parser (Paul Guyot)

The code would try to read one more byte than it should, which would crash in corner cases. Since contrib/test_parser is only example code, this is not a security issue in itself, but bad example code is still bad.

- Use __sync_lock_test_and_set() for spinlocks on ARM, if available (Martin Pitt)

This function replaces our previous use of the SWPB instruction, which is deprecated and not available on ARMv6 and later. Reports suggest that the old code doesn't fail in an obvious way on recent ARM boards, but simply doesn't interlock concurrent accesses, leading to bizarre failures in multiprocess operation.

- Use -fexcess-precision=standard option when building with gcc versions that accept it (Andrew Dunstan)

This prevents assorted scenarios wherein recent versions of gcc will produce creative results.

- Allow use of threaded Python on FreeBSD (Chris Rees)

Our configure script previously believed that this combination wouldn't work; but FreeBSD fixed the problem, so remove that error check.

### E.171. Release 8.4.10

**Release date:** 2011-12-05

This release contains a variety of fixes from 8.4.9. For information about new features in the 8.4 major release, see Section E.181.

#### E.171.1. Migration to Version 8.4.10

A dump/restore is not required for those running 8.4.X.
However, a longstanding error was discovered in the definition of the `information_schema.referential_constraints` view. If you rely on correct results from that view, you should replace its definition as explained in the first changelog item below.

Also, if you are upgrading from a version earlier than 8.4.8, see Section E.173.

**E.171.2. Changes**

- **Fix bugs in `information_schema.referential_constraints` view (Tom Lane)**

  This view was being insufficiently careful about matching the foreign-key constraint to the depended-on primary or unique key constraint. That could result in failure to show a foreign key constraint at all, or showing it multiple times, or claiming that it depends on a different constraint than the one it really does.

  Since the view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can (as a superuser) drop the `information_schema` schema then re-create it by sourcing `SHAREDIR/information_schema.sql`. (Run `pg_config --sharedir` if you're uncertain where `SHAREDIR` is.) This must be repeated in each database to be fixed.

- **Fix incorrect replay of WAL records for GIN index updates (Tom Lane)**

  This could result in transiently failing to find index entries after a crash, or on a hot-standby server. The problem would be repaired by the next `VACUUM` of the index, however.

- **Fix TOAST-related data corruption during `CREATE TABLE dest AS SELECT * FROM src` or `INSERT INTO dest SELECT * FROM src` (Tom Lane)**

  If a table has been modified by `ALTER TABLE ADD COLUMN`, attempts to copy its data verbatim to another table could produce corrupt results in certain corner cases. The problem can only manifest in this precise form in 8.4 and later, but we patched earlier versions as well in case there are other code paths that could trigger the same bug.

- **Fix race condition during toast table access from stale syscache entries (Tom Lane)**

  The typical symptom was transient errors like “missing chunk number 0 for toast value NNNNN in `pg_toast_2619`”, where the cited toast table would always belong to a system catalog.

- **Track dependencies of functions on items used in parameter default expressions (Tom Lane)**

  Previously, a referenced object could be dropped without having dropped or modified the function, leading to misbehavior when the function was used. Note that merely installing this update will not fix the missing dependency entries; to do that, you'd need to `CREATE OR REPLACE` each such function afterwards. If you have functions whose defaults depend on non-built-in objects, doing so is recommended.

- **Allow inlining of set-returning SQL functions with multiple OUT parameters (Tom Lane)**

- **Make `DatumGetInetP()` unpack inet datums that have a 1-byte header, and add a new macro, `DatumGetInetPP()`, that does not (Heikki Linnakangas)**

  This change affects no core code, but might prevent crashes in add-on code that expects `DatumGetInetP()` to produce an unpacked datum as per usual convention.

- **Improve locale support in `money` type's input and output (Tom Lane)**

  Aside from not supporting all standard `lc_monetary` formatting options, the input and output functions were inconsistent, meaning there were locales in which dumped `money` values could not be re-read.
• Don’t let `transform_null_equals` affect `CASE foo WHEN NULL ...` constructs (Heikki Linnakangas)

`transform_null_equals` is only supposed to affect `foo = NULL` expressions written directly by the user, not equality checks generated internally by this form of `CASE`.

• Change foreign-key trigger creation order to better support self-referential foreign keys (Tom Lane)

For a cascading foreign key that references its own table, a row update will fire both the `ON UPDATE` trigger and the `CHECK` trigger as one event. The `ON UPDATE` trigger must execute first, else the `CHECK` will check a non-final state of the row and possibly throw an inappropriate error. However, the firing order of these triggers is determined by their names, which generally sort in creation order since the triggers have auto-generated names following the convention “RI_ConstraintTrigger_NNNN”. A proper fix would require modifying that convention, which we will do in 9.2, but it seems risky to change it in existing releases. So this patch just changes the creation order of the triggers. Users encountering this type of error should drop and re-create the foreign key constraint to get its triggers into the right order.

• Avoid floating-point underflow while tracking buffer allocation rate (Greg Matthews)

While harmless in itself, on certain platforms this would result in annoying kernel log messages.

• Preserve configuration file name and line number values when starting child processes under Windows (Tom Lane)

Formerly, these would not be displayed correctly in the `pg_settings` view.

• Preserve blank lines within commands in `psql`'s command history (Robert Haas)

The former behavior could cause problems if an empty line was removed from within a string literal, for example.

• Fix `pg_dump` to dump user-defined casts between auto-generated types, such as table rowtypes (Tom Lane)

• Use the preferred version of `xsubpp` to build PL/Perl, not necessarily the operating system’s main copy (David Wheeler and Alex Hunsaker)

• Fix incorrect coding in `contrib/dict_int` and `contrib/dict_xsyn` (Tom Lane)

Some functions incorrectly assumed that memory returned by `palloc()` is guaranteed zeroed.

• Honor query cancel interrupts promptly in `pgstatindex()` (Robert Haas)

• Ensure VPATH builds properly install all server header files (Peter Eisentraut)

• Shorten file names reported in verbose error messages (Peter Eisentraut)

Regular builds have always reported just the name of the C file containing the error message call, but VPATH builds formerly reported an absolute path name.

• Fix interpretation of Windows timezone names for Central America (Tom Lane)

Map “Central America Standard Time” to `CST6`, not `CST6CDT`, because DST is generally not observed anywhere in Central America.

• Update time zone data files to tzdata release 2011n for DST law changes in Brazil, Cuba, Fiji, Palestine, Russia, and Samoa; also historical corrections for Alaska and British East Africa.

E.172. Release 8.4.9

Release date: 2011-09-26
This release contains a variety of fixes from 8.4.8. For information about new features in the 8.4 major release, see Section E.181.

E.172.1. Migration to Version 8.4.9

A dump/restore is not required for those running 8.4.X.

However, if you are upgrading from a version earlier than 8.4.8, see Section E.173.

E.172.2. Changes

- Fix bugs in indexing of in-doubt HOT-updated tuples (Tom Lane)
  These bugs could result in index corruption after reindexing a system catalog. They are not believed to affect user indexes.

- Fix multiple bugs in GiST index page split processing (Heikki Linnakangas)
  The probability of occurrence was low, but these could lead to index corruption.

- Fix possible buffer overrun in tsvector_concat() (Tom Lane)
  The function could underestimate the amount of memory needed for its result, leading to server crashes.

- Fix crash in xml_recv when processing a "standalone" parameter (Tom Lane)

- Make pg_options_to_table return NULL for an option with no value (Tom Lane)
  Previously such cases would result in a server crash.

- Avoid possibly accessing off the end of memory in ANALYZE and in SJIS-2004 encoding conversion (Noah Misch)
  This fixes some very-low-probability server crash scenarios.

- Prevent intermittent hang in interactions of startup process with bgwriter process (Simon Riggs)
  This affected recovery in non-hot-standby cases.

- Fix race condition in relcache init file invalidation (Tom Lane)
  There was a window wherein a new backend process could read a stale init file but miss the inval messages that would tell it the data is stale. The result would be bizarre failures in catalog accesses, typically “could not read block 0 in file ...” later during startup.

- Fix memory leak at end of a GiST index scan (Tom Lane)
  Commands that perform many separate GiST index scans, such as verification of a new GiST-based exclusion constraint on a table already containing many rows, could transiently require large amounts of memory due to this leak.

- Fix incorrect memory accounting (leading to possible memory bloat) in tuplestores supporting holdable cursors and plpgsql’s RETURN NEXT command (Tom Lane)

- Fix performance problem when constructing a large, lossy bitmap (Tom Lane)

- Fix join selectivity estimation for unique columns (Tom Lane)
  This fixes an erroneous planner heuristic that could lead to poor estimates of the result size of a join.
Release Notes

- Fix nested PlaceHolderVar expressions that appear only in sub-select target lists (Tom Lane)
  This mistake could result in outputs of an outer join incorrectly appearing as NULL.
- Allow nested EXISTS queries to be optimized properly (Tom Lane)
- Fix array- and path-creating functions to ensure padding bytes are zeroes (Tom Lane)
  This avoids some situations where the planner will think that semantically-equal constants are not equal, resulting in poor optimization.
- Fix EXPLAIN to handle gating Result nodes within inner-indexscan subplans (Tom Lane)
  The usual symptom of this oversight was “bogus varno” errors.
- Work around gcc 4.6.0 bug that breaks WAL replay (Tom Lane)
  This could lead to loss of committed transactions after a server crash.
- Fix dump bug for VALUES in a view (Tom Lane)
- Disallow SELECT FOR UPDATE/SHARE on sequences (Tom Lane)
  This operation doesn't work as expected and can lead to failures.
- Fix VACUUM so that it always updates pg_class.reltuples/relpages (Tom Lane)
  This fixes some scenarios where autovacuum could make increasingly poor decisions about when to vacuum tables.
- Defend against integer overflow when computing size of a hash table (Tom Lane)
- Fix cases where CLUSTER might attempt to access already-removed TOAST data (Tom Lane)
- Fix portability bugs in use of credentials control messages for “peer” authentication (Tom Lane)
  The typical symptom of this problem was “The function requested is not supported” errors during SSPI login.
- Throw an error if pg_hba.conf contains hostssl but SSL is disabled (Tom Lane)
  This was concluded to be more user-friendly than the previous behavior of silently ignoring such lines.
- Fix typo in pg_srand48 seed initialization (Andres Freund)
  This led to failure to use all bits of the provided seed. This function is not used on most platforms (only those without srandom), and the potential security exposure from a less-random-than-expected seed seems minimal in any case.
- Avoid integer overflow when the sum of LIMIT and OFFSET values exceeds 2^63 (Heikki Linnakangas)
- Add overflow checks to int4 and int8 versions of generate_series() (Robert Haas)
- Fix trailing-zero removal in to_char() (Marti Raudsepp)
  In a format with FM and no digit positions after the decimal point, zeroes to the left of the decimal point could be removed incorrectly.
- Fix pg_size_pretty() to avoid overflow for inputs close to 2^63 (Tom Lane)
• Weaken plpgsql’s check for typmod matching in record values (Tom Lane)

An overly enthusiastic check could lead to discarding length modifiers that should have been kept.

• Correctly handle quotes in locale names during initdb (Heikki Linnakangas)

The case can arise with some Windows locales, such as “People’s Republic of China”.

• Fix pg_upgrade to preserve toast tables’ relfrozenxids during an upgrade from 8.3 (Bruce Momjian)

Failure to do this could lead to pg_clog files being removed too soon after the upgrade.

• In pg_ctl, support silent mode for service registrations on Windows (MauMau)

• Fix psql’s counting of script file line numbers during COPY from a different file (Tom Lane)

• Fix pg_restore's direct-to-database mode for standard_conforming_strings (Tom Lane)

pg_restore could emit incorrect commands when restoring directly to a database server from an archive file that had been made with standard_conforming_strings set to on.

• Be more user-friendly about unsupported cases for parallel pg_restore (Tom Lane)

This change ensures that such cases are detected and reported before any restore actions have been taken.

• Fix write-past-buffer-end and memory leak in libpq’s LDAP service lookup code (Albe Laurenz)

• In libpq, avoid failures when using nonblocking I/O and an SSL connection (Martin Pihlak, Tom Lane)

• Improve libpq's handling of failures during connection startup (Tom Lane)

In particular, the response to a server report of fork() failure during SSL connection startup is now saner.

• Improve libpq's error reporting for SSL failures (Tom Lane)

• Fix PQsetvalue() to avoid possible crash when adding a new tuple to a PGresult originally obtained from a server query (Andrew Chernow)

• Make ecpglib write double values with 15 digits precision (Akira Kurosawa)

• In ecpglib, be sure LC_NUMERIC setting is restored after an error (Michael Meskes)

• Apply upstream fix for blowfish signed-character bug (CVE-2011-2483) (Tom Lane)

contrib/pg_crypto’s blowfish encryption code could give wrong results on platforms where char is signed (which is most), leading to encrypted passwords being weaker than they should be.

• Fix memory leak in contrib/seg (Heikki Linnakangas)

• Fix pgstatindex() to give consistent results for empty indexes (Tom Lane)

• Allow building with perl 5.14 (Alex Hunsaker)

• Update configure script's method for probing existence of system functions (Tom Lane)

The version of autoconf we used in 8.3 and 8.2 could be fooled by compilers that perform link-time optimization.

• Fix assorted issues with build and install file paths containing spaces (Tom Lane)
• Update time zone data files to tzdata release 2011i for DST law changes in Canada, Egypt, Russia, Samoa, and South Sudan.

E.173. Release 8.4.8

Release date: 2011-04-18

This release contains a variety of fixes from 8.4.7. For information about new features in the 8.4 major release, see Section E.181.

E.173.1. Migration to Version 8.4.8

A dump/restore is not required for those running 8.4.X.

However, if your installation was upgraded from a previous major release by running pg_upgrade, you should take action to prevent possible data loss due to a now-fixed bug in pg_upgrade. The recommended solution is to run VACUUM FREEZE on all TOAST tables. More information is available at http://wiki.postgresql.org/wiki/20110408pg_upgrade_fix.

Also, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

E.173.2. Changes

• Fix pg_upgrade's handling of TOAST tables (Bruce Momjian)

The pg_class.relrofzenxid value for TOAST tables was not correctly copied into the new installation during pg_upgrade. This could later result in pg_clog files being discarded while they were still needed to validate tuples in the TOAST tables, leading to “could not access status of transaction” failures.

This error poses a significant risk of data loss for installations that have been upgraded with pg_upgrade. This patch corrects the problem for future uses of pg_upgrade, but does not in itself cure the issue in installations that have been processed with a buggy version of pg_upgrade.

• Suppress incorrect “PD_ALL_VISIBLE flag was incorrectly set” warning (Heikki Linnakangas)

VACUUM would sometimes issue this warning in cases that are actually valid.

• Disallow including a composite type in itself (Tom Lane)

This prevents scenarios wherein the server could recurse infinitely while processing the composite type. While there are some possible uses for such a structure, they don't seem compelling enough to justify the effort required to make sure it always works safely.

• Avoid potential deadlock during catalog cache initialization (Nikhil Sontakke)

In some cases the cache loading code would acquire share lock on a system index before locking the index's catalog. This could deadlock against processes trying to acquire exclusive locks in the other, more standard order.

• Fix dangling-pointer problem in BEFORE ROW UPDATE trigger handling when there was a concurrent update to the target tuple (Tom Lane)

This bug has been observed to result in intermittent “cannot extract system attribute from virtual tuple” failures while trying to do UPDATE RETURNING ctid. There is a very small probability of more serious errors, such as generating incorrect index entries for the updated tuple.

• Disallow DROP TABLE when there are pending deferred trigger events for the table (Tom Lane)

---

4 http://wiki.postgresql.org/wiki/20110408pg_upgrade_fix
Formerly the drop would go through, leading to “could not open relation with OID nnn” errors when the triggers were eventually fired.

- Prevent crash triggered by constant-false WHERE conditions during GEQO optimization (Tom Lane)
- Improve planner’s handling of semi-join and anti-join cases (Tom Lane)
- Fix selectivity estimation for text search to account for NULLs (Jesper Krogh)
- Improve PL/pgSQL’s ability to handle row types with dropped columns (Pavel Stehule)
  This is a back-patch of fixes previously made in 9.0.
- Fix PL/Python memory leak involving array slices (Daniel Popowich)
- Fix pg_restore to cope with long lines (over 1KB) in TOC files (Tom Lane)
- Put in more safeguards against crashing due to division-by-zero with overly enthusiastic compiler optimization (Aurelien Jarno)
- Support use of dlopen() in FreeBSD and OpenBSD on MIPS (Tom Lane)
  There was a hard-wired assumption that this system function was not available on MIPS hardware on these systems. Use a compile-time test instead, since more recent versions have it.
- Fix compilation failures on HP-UX (Heikki Linnakangas)
- Fix version-incompatibility problem with libintl on Windows (Hiroshi Inoue)
- Fix usage of xcopy in Windows build scripts to work correctly under Windows 7 (Andrew Dunstan)
  This affects the build scripts only, not installation or usage.
- Fix path separator used by pg_regress on Cygwin (Andrew Dunstan)
- Update time zone data files to tzdata release 2011f for DST law changes in Chile, Cuba, Falkland Islands, Morocco, Samoa, and Turkey; also historical corrections for South Australia, Alaska, and Hawaii.

E.174. Release 8.4.7

Release date: 2011-01-31

This release contains a variety of fixes from 8.4.6. For information about new features in the 8.4 major release, see Section E.181.

E.174.1. Migration to Version 8.4.7

A dump/restore is not required for those running 8.4.X. However, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

E.174.2. Changes

- Avoid failures when EXPLAIN tries to display a simple-form CASE expression (Tom Lane)
  If the CASE’s test expression was a constant, the planner could simplify the CASE into a form that confused the expression-display code, resulting in “unexpected CASE WHEN clause” errors.
• Fix assignment to an array slice that is before the existing range of subscripts (Tom Lane)

If there was a gap between the newly added subscripts and the first pre-existing subscript, the code miscalculated how many entries needed to be copied from the old array's null bitmap, potentially leading to data corruption or crash.

• Avoid unexpected conversion overflow in planner for very distant date values (Tom Lane)

The date type supports a wider range of dates than can be represented by the timestamp types, but the planner assumed it could always convert a date to timestamp with impunity.

• Fix pg_restore's text output for large objects (BLOBs) when standard_conforming_strings is on (Tom Lane)

Although restoring directly to a database worked correctly, string escaping was incorrect if pg_restore was asked for SQL text output and standard_conforming_strings had been enabled in the source database.

• Fix erroneous parsing of tsquery values containing ... & !(subexpression) | ... (Tom Lane)

Queries containing this combination of operators were not executed correctly. The same error existed in contrib/intarray's query_int type and contrib/ltree's ltxtquery type.

• Fix buffer overrun in contrib/intarray's input function for the query_int type (Apple)

This bug is a security risk since the function's return address could be overwritten. Thanks to Apple Inc's security team for reporting this issue and supplying the fix. (CVE-2010-4015)

• Fix bug in contrib/seg's GiST picksplit algorithm (Alexander Korotkov)

This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a seg column. If you have such an index, consider REINDEXing it after installing this update. (This is identical to the bug that was fixed in contrib/cube in the previous update.)

E.175. Release 8.4.6

Release date: 2010-12-16

This release contains a variety of fixes from 8.4.5. For information about new features in the 8.4 major release, see Section E.181.

E.175.1. Migration to Version 8.4.6

A dump/restore is not required for those running 8.4.X. However, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

E.175.2. Changes

• Force the default wal_sync_method to be fdatasync on Linux (Tom Lane, Marti Raudsepp)

The default on Linux has actually been fdatasync for many years, but recent kernel changes caused PostgreSQL to choose open_dataasync instead. This choice did not result in any performance improvement, and caused outright failures on certain filesystems, notably ext4 with the data=journal mount option.

• Fix assorted bugs in WAL replay logic for GIN indexes (Tom Lane)

This could result in “bad buffer id: 0” failures or corruption of index contents during replication.
- Fix recovery from base backup when the starting checkpoint WAL record is not in the same WAL segment as its redo point (Jeff Davis)

- Fix persistent slowdown of autovacuum workers when multiple workers remain active for a long time (Tom Lane)

  The effective vacuum_cost_limit for an autovacuum worker could drop to nearly zero if it processed enough tables, causing it to run extremely slowly.

- Add support for detecting register-stack overrun on IA64 (Tom Lane)

  The IA64 architecture has two hardware stacks. Full prevention of stack-overrun failures requires checking both.

- Add a check for stack overflow in copyObject() (Tom Lane)

  Certain code paths could crash due to stack overflow given a sufficiently complex query.

- Fix detection of page splits in temporary GiST indexes (Heikki Linnakangas)

  It is possible to have a “concurrent” page split in a temporary index, if for example there is an open cursor scanning the index when an insertion is done. GiST failed to detect this case and hence could deliver wrong results when execution of the cursor continued.

- Fix error checking during early connection processing (Tom Lane)

  The check for too many child processes was skipped in some cases, possibly leading to postmaster crash when attempting to add the new child process to fixed-size arrays.

- Improve efficiency of window functions (Tom Lane)

  Certain cases where a large number of tuples needed to be read in advance, but work_mem was large enough to allow them all to be held in memory, were unexpectedly slow. percent_rank(), cume_dist(), and ntile() in particular were subject to this problem.

- Avoid memory leakage while ANALYZE'ing complex index expressions (Tom Lane)

- Ensure an index that uses a whole-row Var still depends on its table (Tom Lane)

  An index declared like create index i on t (foo(t.*) would not automatically get dropped when its table was dropped.

- Do not “inline” a SQL function with multiple OUT parameters (Tom Lane)

  This avoids a possible crash due to loss of information about the expected result rowtype.

- Behave correctly if ORDER BY, LIMIT, FOR UPDATE, or WITH is attached to the VALUES part of INSERT ... VALUES (Tom Lane)

- Fix constant-folding of COALESCE() expressions (Tom Lane)

  The planner would sometimes attempt to evaluate sub-expressions that in fact could never be reached, possibly leading to unexpected errors.

- Fix postmaster crash when connection acceptance (accept() or one of the calls made immediately after it) fails, and the postmaster was compiled with GSSAPI support (Alexander Chernikov)

- Fix missed unlink of temporary files when log_temp_files is active (Tom Lane)

  If an error occurred while attempting to emit the log message, the unlink was not done, resulting in accumulation of temp files.
• Add print functionality for InhRelation nodes (Tom Lane)
  This avoids a failure when debug_print_parse is enabled and certain types of query are executed.
• Fix incorrect calculation of distance from a point to a horizontal line segment (Tom Lane)
  This bug affected several different geometric distance-measurement operators.
• Fix incorrect calculation of transaction status in ecpg (Itagaki Takahiro)
• Fix PL/pgSQL's handling of “simple” expressions to not fail in recursion or error-recovery cases (Tom Lane)
• Fix PL/Python's handling of set-returning functions (Jan Urbanski)
• Fix bug in contrib/cube's GiST picksplit algorithm (Alexander Korotkov)
  This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a cube column. If you have such an index, consider REINDEXing it after installing this update.
• Don't emit “identifier will be truncated” notices in contrib/dblink except when creating new connections (Itagaki Takahiro)
• Fix potential coredump on missing public key in contrib/pgcrypto (Marti Raudsepp)
• Fix memory leak in contrib/xml2's XPath query functions (Tom Lane)
• Update time zone data files to tzdata release 2010o for DST law changes in Fiji and Samoa; also historical corrections for Hong Kong.

E.176. Release 8.4.5

Release date: 2010-10-04

This release contains a variety of fixes from 8.4.4. For information about new features in the 8.4 major release, see Section E.181.

E.176.1. Migration to Version 8.4.5

A dump/restore is not required for those running 8.4.X. However, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

E.176.2. Changes

• Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)
  This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function's owner.
  The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only
one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

- Prevent possible crashes in `pg_get_expr()` by disallowing it from being called with an argument that is not one of the system catalog columns it's intended to be used with (Heikki Linnakangas, Tom Lane)

- Treat exit code 128 (ERROR_WAIT_NO_CHILDREN) as non-fatal on Windows (Magnus Hagander)

Under high load, Windows processes will sometimes fail at startup with this error code. Formerly the postmaster treated this as a panic condition and restarted the whole database, but that seems to be an overreaction.

- Fix incorrect placement of placeholder evaluation (Tom Lane)

This bug could result in query outputs being non-null when they should be null, in cases where the inner side of an outer join is a sub-select with non-strict expressions in its output list.

- Fix possible duplicate scans of `UNION ALL` member relations (Tom Lane)

- Fix "cannot handle unplanned sub-select" error (Tom Lane)

This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.

- Fix mishandling of whole-row Vars that reference a view or sub-select and appear within a nested sub-select (Tom Lane)

- Fix mishandling of cross-type `IN` comparisons (Tom Lane)

This could result in failures if the planner tried to implement an `IN` join with a sort-then-unique-then-plain-join plan.

- Fix computation of `ANALYZE` statistics for `tsvector` columns (Jan Urbanski)

The original coding could produce incorrect statistics, leading to poor plan choices later.

- Improve planner’s estimate of memory used by `array_agg()`, `string_agg()`, and similar aggregate functions (Hitoshi Harada)

The previous drastic underestimate could lead to out-of-memory failures due to inappropriate choice of a hash-aggregation plan.

- Fix failure to mark cached plans as transient (Tom Lane)

If a plan is prepared while `CREATE INDEX CONCURRENTLY` is in progress for one of the referenced tables, it is supposed to be re-planned once the index is ready for use. This was not happening reliably.

- Reduce `PANIC` to `ERROR` in some occasionally-reported btree failure cases, and provide additional detail in the resulting error messages (Tom Lane)

This should improve the system's robustness with corrupted indexes.

- Fix incorrect search logic for partial-match queries with GIN indexes (Tom Lane)
Cases involving AND/OR combination of several GIN index conditions didn't always give the right answer, and were sometimes much slower than necessary.

- Prevent show_session_authorization() from crashing within autovacuum processes (Tom Lane)
- Defend against functions returning setof record where not all the returned rows are actually of the same rowtype (Tom Lane)
- Fix possible corruption of pending trigger event lists during subtransaction rollback (Tom Lane)
  This could lead to a crash or incorrect firing of triggers.
- Fix possible failure when hashing a pass-by-reference function result (Tao Ma, Tom Lane)
- Improve merge join's handling of NULLs in the join columns (Tom Lane)
  A merge join can now stop entirely upon reaching the first NULL, if the sort order is such that NULLs sort high.
- Take care to fsync the contents of lockfiles (both postmaster.pid and the socket lockfile) while writing them (Tom Lane)
  This omission could result in corrupted lockfile contents if the machine crashes shortly after post-master start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.
- Avoid recursion while assigning XIDs to heavily-nested subtransactions (Andres Freund, Robert Haas)
  The original coding could result in a crash if there was limited stack space.
- Avoid holding open old WAL segments in the walwriter process (Magnus Hagander, Heikki Linnakangas)
  The previous coding would prevent removal of no-longer-needed segments.
- Fix log_line_prefix's %i escape, which could produce junk early in backend startup (Tom Lane)
- Prevent misinterpretation of partially-specified relation options for TOAST tables (Itagaki Takahiro)
  In particular, fillfactor would be read as zero if any other reloption had been set for the table, leading to serious bloat.
- Fix inheritance count tracking in ALTER TABLE ... ADD CONSTRAINT (Robert Haas)
- Fix possible data corruption in ALTER TABLE ... SET TABLESPACE when archiving is enabled (Jeff Davis)
- Allow CREATE DATABASE and ALTER DATABASE ... SET TABLESPACE to be interrupted by query-cancel (Guillaume Lelarge)
- Improve CREATE INDEX's checking of whether proposed index expressions are immutable (Tom Lane)
- Fix REASSIGN OWNED to handle operator classes and families (Asko Tiidumaa)
- Fix possible core dump when comparing two empty tsquery values (Tom Lane)
- Fix LIKE's handling of patterns containing % followed by _ (Tom Lane)
We've fixed this before, but there were still some incorrectly-handled cases.

• Re-allow input of Julian dates prior to 0001-01-01 AD (Tom Lane)
  Input such as 'J100000'::date worked before 8.4, but was unintentionally broken by added error-checking.

• Fix PL/pgSQL to throw an error, not crash, if a cursor is closed within a FOR loop that is iterating over that cursor (Heikki Linnakangas)

• In PL/Python, defend against null pointer results from PyCObject_AsVoidPtr and PyCObject_FromVoidPtr (Peter Eisentraut)

• In libpq, fix full SSL certificate verification for the case where both host and hostaddr are specified (Tom Lane)

• Make psql recognize DISCARD ALL as a command that should not be encased in a transaction block in autocommit-off mode (Itagaki Takahiro)

• Fix some issues in pg_dump's handling of SQL/MED objects (Tom Lane)
  Notably, pg_dump would always fail if run by a non-superuser, which was not intended.

• Improve pg_dump and pg_restore's handling of non-seekable archive files (Tom Lane, Robert Haas)
  This is important for proper functioning of parallel restore.

• Improve parallel pg_restore's ability to cope with selective restore (-L option) (Tom Lane)
  The original code tended to fail if the -L file commanded a non-default restore ordering.

• Fix ecpg to process data from RETURNING clauses correctly (Michael Meskes)

• Fix some memory leaks in ecpg (Zoltan Boszormenyi)

• Improve contrib/dblink's handling of tables containing dropped columns (Tom Lane)

• Fix connection leak after “duplicate connection name” errors in contrib/dblink (Itagaki Takahiro)

• Fix contrib/dblink to handle connection names longer than 62 bytes correctly (Itagaki Takahiro)

• Add hstore(text, text) function to contrib/hstore (Robert Haas)
  This function is the recommended substitute for the now-deprecated => operator. It was back-patched so that future-proofed code can be used with older server versions. Note that the patch will be effective only after contrib/hstore is installed or reinstalled in a particular database. Users might prefer to execute the CREATE FUNCTION command by hand, instead.

• Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)

• Update time zone data files to tzdata release 2010l for DST law changes in Egypt and Palestine; also historical corrections for Finland.
  This change also adds new names for two Micronesian timezones: Pacific/Chuuk is now preferred over Pacific/Truk (and the preferred abbreviation is CHUT not TRUT) and Pacific/Pohnpei is preferred over Pacific/Ponape.

• Make Windows' "N. Central Asia Standard Time" timezone map to Asia/Novosibirsk, not Asia/Almaty (Magnus Hagander)
Microsoft changed the DST behavior of this zone in the timezone update from KB976098. Asia/Novosibirsk is a better match to its new behavior.

E.177. Release 8.4.4

Release date: 2010-05-17

This release contains a variety of fixes from 8.4.3. For information about new features in the 8.4 major release, see Section E.181.

E.177.1. Migration to Version 8.4.4

A dump/restore is not required for those running 8.4.X. However, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

E.177.2. Changes

• Enforce restrictions in `plperl` using an opmask applied to the whole interpreter, instead of using `Safe.pm` (Tim Bunce, Andrew Dunstan)

Recent developments have convinced us that `Safe.pm` is too insecure to rely on for making `plperl` trustworthy. This change removes use of `Safe.pm` altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include that it is now possible to use Perl's `strict` pragma in a natural way in `plperl`, and that Perl's `$a` and `$b` variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

• Prevent PL/Tcl from executing untrustworthy code from `pltcl_modules` (Tom)

PL/Tcl's feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless `pltcl_modules` is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted "normal" Tcl interpreter unless we are really going to execute a `pltclu` function. (CVE-2010-1170)

• Fix data corruption during WAL replay of `ALTER ... SET TABLESPACE` (Tom)

When `archive_mode` is on, `ALTER ... SET TABLESPACE` generates a WAL record whose replay logic was incorrect. It could write the data to the wrong place, leading to possibly-unrecoverable data corruption. Data corruption would be observed on standby slaves, and could occur on the master as well if a database crash and recovery occurred after committing the `ALTER` and before the next checkpoint.

• Fix possible crash if a cache reset message is received during rebuild of a relcache entry (Heikki)

This error was introduced in 8.4.3 while fixing a related failure.

• Apply per-function GUC settings while running the language validator for the function (Itagaki Takahiro)

This avoids failures if the function's code is invalid without the setting; an example is that SQL functions may not parse if the `search_path` is not correct.

• Do constraint exclusion for inherited `UPDATE` and `DELETE` target tables when `constraint_exclusion = partition` (Tom)
Due to an oversight, this setting previously only caused constraint exclusion to be checked in
SELECT commands.

- Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

  Previously, if an unprivileged user ran ALTER USER ... RESET ALL for himself, or ALTER
  DATABASE ... RESET ALL for a database he owns, this would remove all special parameter
  settings for the user or database, even ones that are only supposed to be changeable by a superuser.
  Now, the ALTER will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a CONTEXT addition
  would be made to log entries (Tom)

  In some cases the context-printing function would fail because the current transaction had already
  been rolled back when it came time to print a log message.

- Fix erroneous handling of %r parameter in recovery_end_command (Heikki)

  The value always came out zero.

- Ensure the archiver process responds to changes in archive_command as soon as possible (Tom)

- Fix PL/pgSQL's CASE statement to not fail when the case expression is a query that returns no
  rows (Tom)

- Update PL/Perl's ppport.h for modern Perl versions (Andrew)

- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)

- Handle empty-string connect parameters properly in ecpg (Michael)

- Prevent infinite recursion in psql when expanding a variable that refers to itself (Tom)

- Fix psql's \copy to not add spaces around a dot within \copy (select ...) (Tom)

  Addition of spaces around the decimal point in a numeric literal would result in a syntax error.

- Avoid formatting failure in psql when running in a locale context that doesn't match the client_encoding (Tom)

- Fix unnecessary “GIN indexes do not support whole-index scans” errors for unsatisfiable queries
  using contrib/intarray operators (Tom)

- Ensure that contrib/pgstattuple functions respond to cancel interrupts promptly (Tatsuhito
  Kasahara)

- Make server startup deal properly with the case that shmget() returns EINVAL for an existing
  shared memory segment (Tom)

  This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entire-
  ly-misleading startup failure complaining that the shared memory request size was too large.

- Avoid possible crashes in syslogger process on Windows (Heikki)

- Deal more robustly with incomplete time zone information in the Windows registry (Magnus)

- Update the set of known Windows time zone names (Magnus)

- Update time zone data files to tzdata release 2010j for DST law changes in Argentina, Australian
  Antarctic, Bangladesh, Mexico, Morocco, Pakistan, Palestine, Russia, Syria, Tunisia; also historical
  corrections for Taiwan.
Also, add PKST (Pakistan Summer Time) to the default set of timezone abbreviations.

**E.178. Release 8.4.3**

**Release date:** 2010-03-15

This release contains a variety of fixes from 8.4.2. For information about new features in the 8.4 major release, see Section E.181.

**E.178.1. Migration to Version 8.4.3**

A dump/restore is not required for those running 8.4.X. However, if you are upgrading from a version earlier than 8.4.2, see Section E.179.

**E.178.2. Changes**

- Add new configuration parameter `ssl_renegotiation_limit` to control how often we do session key renegotiation for an SSL connection (Magnus)

  This can be set to zero to disable renegotiation completely, which may be required if a broken SSL library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that cause renegotiation attempts to fail.

- Fix possible deadlock during backend startup (Tom)

- Fix possible crashes due to not handling errors during relcache reload cleanly (Tom)

- Fix possible crash due to use of dangling pointer to a cached plan (Tatsuo)

- Fix possible crash due to overenthusiastic invalidation of cached plan for ROLLBACK (Tom)

- Fix possible crashes when trying to recover from a failure in subtransaction start (Tom)

- Fix server memory leak associated with use of savepoints and a client encoding different from server's encoding (Tom)

- Fix incorrect WAL data emitted during end-of-recovery cleanup of a GIST index page split (Yoichi Hirai)

  This would result in index corruption, or even more likely an error during WAL replay, if we were unlucky enough to crash during end-of-recovery cleanup after having completed an incomplete GIST insertion.

- Fix bug in WAL redo cleanup method for GIN indexes (Heikki)

- Fix incorrect comparison of scan key in GIN index search (Teodor)

- Make `substring()` for bit types treat any negative length as meaning “all the rest of the string” (Tom)

  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

- Fix integer-to-bit-string conversions to handle the first fractional byte correctly when the output bit width is wider than the given integer by something other than a multiple of 8 bits (Tom)

- Fix some cases of pathologically slow regular expression matching (Tom)

- Fix bug occurring when trying to inline a SQL function that returns a set of a composite type that contains dropped columns (Tom)
• Fix bug with trying to update a field of an element of a composite-type array column (Tom)

• Avoid failure when EXPLAIN has to print a FieldStore or assignment ArrayRef expression (Tom)

These cases can arise now that EXPLAIN VERbose tries to print plan node target lists.

• Avoid an unnecessary coercion failure in some cases where an undecorated literal string appears in a subquery within UNION/INTERSECT/EXCEPT (Tom)

This fixes a regression for some cases that worked before 8.4.

• Avoid undesirable rowtype compatibility check failures in some cases where a whole-row Var has a rowtype that contains dropped columns (Tom)

• Fix the STOP WAL LOCATION entry in backup history files to report the next WAL segment's name when the end location is exactly at a segment boundary (Itagaki Takahiro)

• Always pass the catalog ID to an option validator function specified in CREATE FOREIGN DATA WRAPPER (Martin Pihlak)

• Fix some more cases of temporary-file leakage (Heikki)

This corrects a problem introduced in the previous minor release. One case that failed is when a plpgsql function returning set is called within another function's exception handler.

• Add support for doing FULL JOIN ON FALSE (Tom)

This prevents a regression from pre-8.4 releases for some queries that can now be simplified to a constant-false join condition.

• Improve constraint exclusion processing of boolean-variable cases, in particular make it possible to exclude a partition that has a “bool_column = false” constraint (Tom)

• Prevent treating an INOUT cast as representing binary compatibility (Heikki)

• Include column name in the message when warning about inability to grant or revoke column-level privileges (Stephen Frost)

This is more useful than before and helps to prevent confusion when a REVOKE generates multiple messages, which formerly appeared to be duplicates.

• When reading pg_hba.conf and related files, do not treat @something as a file inclusion request if the @ appears inside quote marks; also, never treat @ by itself as a file inclusion request (Tom)

This prevents erratic behavior if a role or database name starts with @. If you need to include a file whose path name contains spaces, you can still do so, but you must write @"/path to/file" rather than putting the quotes around the whole construct.

• Prevent infinite loop on some platforms if a directory is named as an inclusion target in pg_hba.conf and related files (Tom)

• Fix possible infinite loop if SSL_read or SSL_write fails without setting errno (Tom)

This is reportedly possible with some Windows versions of OpenSSL.

• Disallow GSSAPI authentication on local connections, since it requires a hostname to function correctly (Magnus)

• Protect ecpg against applications freeing strings unexpectedly (Michael)

• Make ecpg report the proper SQLSTATE if the connection disappears (Michael)
Release Notes

- Fix translation of cell contents in psql `\d` output (Heikki)
- Fix psql's `numericlocale` option to not format strings it shouldn't in latex and troff output formats (Heikki)
- Fix a small per-query memory leak in psql (Tom)
- Make psql return the correct exit status (3) when `ON_ERROR_STOP` and `--single-transaction` are both specified and an error occurs during the implied COMMIT (Bruce)
- Fix `pg_dump`'s output of permissions for foreign servers (Heikki)
- Fix possible crash in parallel `pg_restore` due to out-of-range dependency IDs (Tom)
- Fix plpgsql failure in one case where a composite column is set to NULL (Tom)
- Fix possible failure when calling PL/Perl functions from PL/PerlU or vice versa (Tim Bunce)
- Add `volatile` markings in PL/Python to avoid possible compiler-specific misbehavior (Zdenek Kotala)
- Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)
  The only known symptom of this oversight is that the Tcl `clock` command misbehaves if using Tcl 8.5 or later.
- Prevent `ExecutorEnd` from being run on portals created within a failed transaction or subtransaction (Tom)
  This is known to cause issues when using `contrib/auto_explain`.
- Prevent crash in `contrib/dblink` when too many key columns are specified to a `dblink_build_sql_*` function (Rushabh Lathia, Joe Conway)
- Allow zero-dimensional arrays in `contrib/ltree` operations (Tom)
  This case was formerly rejected as an error, but it's more convenient to treat it the same as a zero-element array. In particular this avoids unnecessary failures when an `ltree` operation is applied to the result of `ARRAY(SELECT ...) and the sub-select returns no rows.
- Fix assorted crashes in `contrib/xml2` caused by sloppy memory management (Tom)
- Make building of `contrib/xml2` more robust on Windows (Andrew)
- Fix race condition in Windows signal handling (Radu Ilie)
  One known symptom of this bug is that rows in `pg_listener` could be dropped under heavy load.
- Make the configure script report failure if the C compiler does not provide a working 64-bit integer datatype (Tom)
  This case has been broken for some time, and no longer seems worth supporting, so just reject it at configure time instead.
- Update time zone data files to tzdata release 2010e for DST law changes in Bangladesh, Chile, Fiji, Mexico, Paraguay, Samoa.

E.179. Release 8.4.2

Release date: 2009-12-14

This release contains a variety of fixes from 8.4.1. For information about new features in the 8.4 major release, see Section E.181.
E.179.1. Migration to Version 8.4.2

A dump/restore is not required for those running 8.4.X. However, if you have any hash indexes, you should REINDEX them after updating to 8.4.2, to repair possible damage.

E.179.2. Changes

• Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)

  This change prevents allegedly-immutable index functions from possibly subverting a superuser’s session (CVE-2009-4136).

• Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)

  This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).

• Fix hash index corruption (Tom)

  The 8.4 change that made hash indexes keep entries sorted by hash value failed to update the bucket splitting and compaction routines to preserve the ordering. So application of either of those operations could lead to permanent corruption of an index, in the sense that searches might fail to find entries that are present. To deal with this, it is recommended to REINDEX any hash indexes you may have after installing this update.

• Fix possible crash during backend-startup-time cache initialization (Tom)

• Avoid crash on empty thesaurus dictionary (Tom)

• Prevent signals from interrupting VACUUM at unsafe times (Alvaro)

  This fix prevents a PANIC if a VACUUM FULL is canceled after it’s already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

• Fix possible crash due to integer overflow in hash table size calculation (Tom)

  This could occur with extremely large planner estimates for the size of a hashjoin’s result.

• Fix crash if a DROP is attempted on an internally-dependent object (Tom)

• Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)

• Ensure that shared tuple-level locks held by prepared transactions are not ignored (Heikki)

• Fix premature drop of temporary files used for a cursor that is accessed within a subtransaction (Heikki)

• Fix memory leak in syslogger process when rotating to a new CSV logfile (Tom)

• Fix memory leak in postmaster when re-parsing pg_hba.conf (Tom)

• Fix Windows permission-downgrade logic (Jesse Morris)

  This fixes some cases where the database failed to start on Windows, often with misleading error messages such as “could not locate matching postgres executable”.

• Make FOR UPDATE/SHARE in the primary query not propagate into WITH queries (Tom)

  For example, in
WITH w AS (SELECT * FROM foo) SELECT * FROM w, bar ... FOR UPDATE

the FOR UPDATE will now affect bar but not foo. This is more useful and consistent than the original 8.4 behavior, which tried to propagate FOR UPDATE into the WITH query but always failed due to assorted implementation restrictions. It also follows the design rule that WITH queries are executed as if independent of the main query.

• Fix bug with a WITH RECURSIVE query immediately inside another one (Tom)
• Fix concurrency bug in hash indexes (Tom)
  Concurrent insertions could cause index scans to transiently report wrong results.
• Fix incorrect logic for GiST index page splits, when the split depends on a non-first column of the index (Paul Ramsey)
• Fix wrong search results for a multi-column GIN index with fastupdate enabled (Teodor)
• Fix bugs in WAL entry creation for GIN indexes (Tom)
  These bugs were masked when full_page_writes was on, but with it off a WAL replay failure was certain if a crash occurred before the next checkpoint.
• Don’t error out if recycling or removing an old WAL file fails at the end of checkpoint (Heikki)
  It’s better to treat the problem as non-fatal and allow the checkpoint to complete. Future checkpoints will retry the removal. Such problems are not expected in normal operation, but have been seen to be caused by misdesigned Windows anti-virus and backup software.
• Ensure WAL files aren’t repeatedly archived on Windows (Heikki)
  This is another symptom that could happen if some other process interfered with deletion of a no-longer-needed file.
• Fix PAM password processing to be more robust (Tom)
  The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.
• Raise the maximum authentication token (Kerberos ticket) size in GSSAPI and SSPI authentication methods (Ian Turner)
  While the old 2000-byte limit was more than enough for Unix Kerberos implementations, tickets issued by Windows Domain Controllers can be much larger.
• Ensure that domain constraints are enforced in constructs like ARRAY[...]:domain, where the domain is over an array type (Heikki)
• Fix foreign-key logic for some cases involving composite-type columns as foreign keys (Tom)
• Ensure that a cursor’s snapshot is not modified after it is created (Alvaro)
  This could lead to a cursor delivering wrong results if later operations in the same transaction modify the data the cursor is supposed to return.
• Fix CREATE TABLE to properly merge default expressions coming from different inheritance parent tables (Tom)
  This used to work but was broken in 8.4.
• Re-enable collection of access statistics for sequences (Akira Kurosawa)

This used to work but was broken in 8.3.

• Fix processing of ownership dependencies during CREATE OR REPLACE FUNCTION (Tom)

• Fix incorrect handling of WHERE x=x conditions (Tom)

In some cases these could get ignored as redundant, but they aren’t — they’re equivalent to x IS NOT NULL.

• Fix incorrect plan construction when using hash aggregation to implement DISTINCT for textually identical volatile expressions (Tom)

• Fix Assert failure for a volatile SELECT DISTINCT ON expression (Tom)

• Fix ts_stat() to not fail on an empty tsvector value (Tom)

• Make text search parser accept underscores in XML attributes (Peter)

• Fix encoding handling in xml binary input (Heikki)

If the XML header doesn’t specify an encoding, we now assume UTF-8 by default; the previous handling was inconsistent.

• Fix bug with calling plperl from plperlu or vice versa (Tom)

An error exit from the inner function could result in crashes due to failure to re-select the correct Perl interpreter for the outer function.

• Fix session-lifespan memory leak when a PL/Perl function is redefined (Tom)

• Ensure that Perl arrays are properly converted to PostgreSQL arrays when returned by a set-returning PL/Perl function (Andrew Dunstan, Abhijit Menon-Sen)

This worked correctly already for non-set-returning functions.

• Fix rare crash in exception processing in PL/Python (Peter)

• Fix ecpg problem with comments in DECLARE CURSOR statements (Michael)

• Fix ecpg to not treat recently-added keywords as reserved words (Tom)

This affected the keywords CALLED, CATALOG, DEFINER, ENUM, FOLLOWING, INVOKER, OPTIONS, PARTITION, PRECEDING, RANGE, SECURITY, SERVER, UNBOUNDED, and WRAPPER.

• Re-allow regular expression special characters in psql’s \df function name parameter (Tom)

• In contrib/fuzzystrmatch, correct the calculation of levenshtein distances with non-default costs (Marcin Mank)

• In contrib/pg_standby, disable triggering failover with a signal on Windows (Fujii Masao)

This never did anything useful, because Windows doesn’t have Unix-style signals, but recent changes made it actually crash.

• Put FREEZE and VERBOSE options in the right order in the VACUUM command that contrib/vacuumdb produces (Heikki)

• Fix possible leak of connections when contrib/dblink encounters an error (Tatsuhito Kasahara)
• Ensure psql's flex module is compiled with the correct system header definitions (Tom)

This fixes build failures on platforms where --enable-largefile causes incompatible changes in the generated code.

• Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

• Update the timezone abbreviation files to match current reality (Joachim Wieland)

This includes adding IDT to the default timezone abbreviation set.

• Update time zone data files to tzdata release 2009s for DST law changes in Antarctica, Argentina, Bangladesh, Fiji, Novokuznetsk, Pakistan, Palestine, Samoa, Syria; also historical corrections for Hong Kong.

E.180. Release 8.4.1

Release date: 2009-09-09

This release contains a variety of fixes from 8.4. For information about new features in the 8.4 major release, see Section E.181.

E.180.1. Migration to Version 8.4.1

A dump/restore is not required for those running 8.4.X.

E.180.2. Changes

• Fix WAL page header initialization at the end of archive recovery (Heikki)

This could lead to failure to process the WAL in a subsequent archive recovery.

• Fix “cannot make new WAL entries during recovery” error (Tom)

• Fix problem that could make expired rows visible after a crash (Tom)

This bug involved a page status bit potentially not being set correctly after a server crash.

• Disallow RESET ROLE and RESET SESSION AUTHORIZATION inside security-definer functions (Tom, Heikki)

This covers a case that was missed in the previous patch that disallowed SET ROLE and SET SESSION AUTHORIZATION inside security-definer functions. (See CVE-2007-6600)

• Make LOAD of an already-loaded loadable module into a no-op (Tom)

Formerly, LOAD would attempt to unload and re-load the module, but this is unsafe and not all that useful.

• Make window function PARTITION BY and ORDER BY items always be interpreted as simple expressions (Tom)

In 8.4.0 these lists were parsed following the rules used for top-level GROUP BY and ORDER BY lists. But this was not correct per the SQL standard, and it led to possible circularity.

• Fix several errors in planning of semi-joins (Tom)

These led to wrong query results in some cases where IN or EXISTS was used together with another join.
• Fix handling of whole-row references to subqueries that are within an outer join (Tom)

An example is `SELECT COUNT(ss.*) FROM ... LEFT JOIN (SELECT ...) ss ON ...`. Here, `ss.*` would be treated as `ROW(NULL,NULL, ...)` for null-extended join rows, which is not the same as a simple NULL. Now it is treated as a simple NULL.

• Fix Windows shared-memory allocation code (Tsutomu Yamada, Magnus)

This bug led to the often-reported “could not reattach to shared memory” error message.

• Fix locale handling with plperl (Heikki)

This bug could cause the server's locale setting to change when a plperl function is called, leading to data corruption.

• Fix handling of reloptions to ensure setting one option doesn’t force default values for others (Itagaki Takahiro)

• Ensure that a “fast shutdown” request will forcibly terminate open sessions, even if a “smart shutdown” was already in progress (Fuji Masao)

• Avoid memory leak for `array_agg()` in GROUP BY queries (Tom)

• Treat `to_char(..., 'TH')` as an uppercase ordinal suffix with 'HH'/'HH12' (Heikki)

   It was previously handled as 'th' (lowercase).

• Include the fractional part in the result of `EXTRACT(second)` and `EXTRACT(milliseconds)` for `time` and `time with time zone` inputs (Tom)

   This has always worked for floating-point datetime configurations, but was broken in the integer datetime code.

• Fix overflow for `INTERVAL 'x ms'` when `x` is more than 2 million and integer datetimes are in use (Alex Hunsaker)

• Improve performance when processing toasted values in index scans (Tom)

   This is particularly useful for PostGIS\(^5\).

• Fix a typo that disabled `commit_delay` (Jeff Janes)

• Output early-startup messages to `postmaster.log` if the server is started in silent mode (Tom)

   Previously such error messages were discarded, leading to difficulty in debugging.

• Remove translated FAQs (Peter)

   They are now on the wiki\(^6\). The main FAQ was moved to the wiki some time ago.

• Fix `pg_ctl` to not go into an infinite loop if `postgresql.conf` is empty (Jeff Davis)

• Fix several errors in `pg_dump`'s `--binary-upgrade` mode (Bruce, Tom)

   `pg_dump` `--binary-upgrade` is used by `pg_migrator`.

• Fix `contrib/xml2`'s `xslt_process()` to properly handle the maximum number of parameters (twenty) (Tom)

• Improve robustness of `libpq`'s code to recover from errors during `COPY FROM STDIN` (Tom)

\(^5\) http://postgis.net/
\(^6\) http://wiki.postgresql.org/wiki/FAQ
• Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)
• Work around gcc bug that causes “floating-point exception” instead of “division by zero” on some platforms (Tom)
• Update time zone data files to tzdata release 2009l for DST law changes in Bangladesh, Egypt, Mauritius.

E.181. Release 8.4

Release date: 2009-07-01

E.181.1. Overview

After many years of development, PostgreSQL has become feature-complete in many areas. This release shows a targeted approach to adding features (e.g., authentication, monitoring, space reuse), and adds capabilities defined in the later SQL standards. The major areas of enhancement are:

• Windowing Functions
• Common Table Expressions and Recursive Queries
• Default and variadic parameters for functions
• Parallel Restore
• Column Permissions
• Per-database locale settings
• Improved hash indexes
• Improved join performance for EXISTS and NOT EXISTS queries
• Easier-to-use Warm Standby
• Automatic sizing of the Free Space Map
• Visibility Map (greatly reduces vacuum overhead for slowly-changing tables)
• Version-aware psql (backslash commands work against older servers)
• Support SSL certificates for user authentication
• Per-function runtime statistics
• Easy editing of functions in psql
• New contrib modules: pg_stat_statements, auto_explain, citext, btree_gin

The above items are explained in more detail in the sections below.

E.181.2. Migration to Version 8.4

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

E.181.2.1. General

• Use 64-bit integer datetimes by default (Neil Conway)
Previously this was selected by configure's --enable-integer-datetimes option. To retain the old behavior, build with --disable-integer-datetimes.

• Remove ipcclean utility command (Bruce)

The utility only worked on a few platforms. Users should use their operating system tools instead.

### E.181.2.2. Server Settings

- Change default setting for `log_min_messages` to `warning` (previously it was `notice`) to reduce log file volume (Tom)

- Change default setting for `max_prepared_transactions` to zero (previously it was 5) (Tom)

- Make `debug_print_parse`, `debug_print_rewritten`, and `debug_print_plan` output appear at `LOG` message level, not `DEBUG1` as formerly (Tom)

- Make `debug_pretty_print` default to `on` (Tom)

- Remove `explain_pretty_print` parameter (no longer needed) (Tom)

- Make `log_temp_files` settable by superusers only, like other logging options (Simon Riggs)

- Remove automatic appending of the epoch timestamp when no `%` escapes are present in `log_filename` (Robert Haas)

  This change was made because some users wanted a fixed log filename, for use with an external log rotation tool.

- Remove `log_restartpoints` from `recovery.conf`; instead use `log_checkpoints` (Simon)

- Remove `krb_realm` and `krb_server_hostname`; these are now set in `pg_hba.conf` instead (Magnus)

- There are also significant changes in `pg_hba.conf`, as described below.

### E.181.2.3. Queries

- Change `TRUNCATE` and `LOCK` to apply to child tables of the specified table(s) (Peter)

  These commands now accept an `ONLY` option that prevents processing child tables; this option must be used if the old behavior is needed.

- `SELECT DISTINCT` and `UNION/INTERSECT/EXCEPT` no longer always produce sorted output (Tom)

  Previously, these types of queries always removed duplicate rows by means of Sort/Unique processing (i.e., sort then remove adjacent duplicates). Now they can be implemented by hashing, which will not produce sorted output. If an application relied on the output being in sorted order, the recommended fix is to add an `ORDER BY` clause. As a short-term workaround, the previous behavior can be restored by disabling `enable_hashagg`, but that is a very performance-expensive fix. `SELECT DISTINCT ON` never uses hashing, however, so its behavior is unchanged.

- Force child tables to inherit `CHECK` constraints from parents (Alex Hunsaker, Nikhil Sontakke, Tom)

  Formerly it was possible to drop such a constraint from a child table, allowing rows that violate the constraint to be visible when scanning the parent table. This was deemed inconsistent, as well as contrary to SQL standard.
• Disallow negative LIMIT or OFFSET values, rather than treating them as zero (Simon)

• Disallow LOCK TABLE outside a transaction block (Tom)

  Such an operation is useless because the lock would be released immediately.

• Sequences now contain an additional start_value column (Zoltan Boszormenyi)

  This supports ALTER SEQUENCE ... RESTART.

E.181.2.4. Functions and Operators

• Make numeric zero raised to a fractional power return 0, rather than throwing an error, and make numeric zero raised to the zero power return 1, rather than error (Bruce)

  This matches the longstanding float8 behavior.

• Allow unary minus of floating-point values to produce minus zero (Tom)

  The changed behavior is more IEEE-standard compliant.

• Throw an error if an escape character is the last character in a LIKE pattern (i.e., it has nothing to escape) (Tom)

  Previously, such an escape character was silently ignored, thus possibly masking application logic errors.

• Remove ~=~ and ~<>~ operators formerly used for LIKE index comparisons (Tom)

  Pattern indexes now use the regular equality operator.

• xpath() now passes its arguments to libxml without any changes (Andrew)

  This means that the XML argument must be a well-formed XML document. The previous coding attempted to allow XML fragments, but it did not work well.

• Make xmlelement() format attribute values just like content values (Peter)

  Previously, attribute values were formatted according to the normal SQL output behavior, which is sometimes at odds with XML rules.

• Rewrite memory management for libxml-using functions (Tom)

  This change should avoid some compatibility problems with use of libxml in PL/Perl and other add-on code.

• Adopt a faster algorithm for hash functions (Kenneth Marshall, based on work of Bob Jenkins)

  Many of the built-in hash functions now deliver different results on little-endian and big-endian platforms.

E.181.2.4.1. Temporal Functions and Operators

• DateStyle no longer controls interval output formatting; instead there is a new variable IntervalStyle (Ron Mayer)

• Improve consistency of handling of fractional seconds in timestamp and interval output (Ron Mayer)

  This may result in displaying a different number of fractional digits than before, or rounding instead of truncating.
• Make to_char()'s localized month/day names depend on LC_TIME, not LC_MESSAGES (Euler Taveira de Oliveira)

• Cause to_date() and to_timestamp() to more consistently report errors for invalid input (Brendan Jurd)

Previous versions would often ignore or silently misread input that did not match the format string. Such cases will now result in an error.

• Fix to_timestamp() to not require upper/lower case matching for meridian (AM/PM) and era (BC/AD) format designations (Brendan Jurd)

For example, input value ad now matches the format string AD.

E.181.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 8.4 and the previous major release.

E.181.3.1. Performance

• Improve optimizer statistics calculations (Jan Urbanski, Tom)

  In particular, estimates for full-text-search operators are greatly improved.

• Allow SELECT DISTINCT and UNION/INTERSECT/EXCEPT to use hashing (Tom)

  This means that these types of queries no longer automatically produce sorted output.

• Create explicit concepts of semi-joins and anti-joins (Tom)

  This work formalizes our previous ad-hoc treatment of IN (SELECT ...) clauses, and extends it to EXISTS and NOT EXISTS clauses. It should result in significantly better planning of EXISTS and NOT EXISTS queries. In general, logically equivalent IN and EXISTS clauses should now have similar performance, whereas previously IN often won.

• Improve optimization of sub-selects beneath outer joins (Tom)

  Formerly, a sub-select or view could not be optimized very well if it appeared within the nullable side of an outer join and contained non-strict expressions (for instance, constants) in its result list.

• Improve the performance of text_position() and related functions by using Boyer-Moore-Horspool searching (David Rowley)

  This is particularly helpful for long search patterns.

• Reduce I/O load of writing the statistics collection file by writing the file only when requested (Martin Pihlak)

• Improve performance for bulk inserts (Robert Haas, Simon)

• Increase the default value of default_statistics_target from 10 to 100 (Greg Sabino Mullane, Tom)

  The maximum value was also increased from 1000 to 10000.

• Perform constraint_exclusion checking by default in queries involving inheritance or UNION ALL (Tom)

  A new constraint_exclusion setting, partition, was added to specify this behavior.
• Allow I/O read-ahead for bitmap index scans (Greg Stark)

The amount of read-ahead is controlled by `effective_io_concurrency`. This feature is available only if the kernel has `posix_fadvise()` support.

• Inline simple set-returning SQL functions in `FROM` clauses (Richard Rowell)

• Improve performance of multi-batch hash joins by providing a special case for join key values that are especially common in the outer relation (Bryce Cutt, Ramon Lawrence)

• Reduce volume of temporary data in multi-batch hash joins by suppressing “physical tlist” optimization (Michael Henderson, Ramon Lawrence)

• Avoid waiting for idle-in-transaction sessions during `CREATE INDEX CONCURRENTLY` (Simon)

• Improve performance of shared cache invalidation (Tom)

E.181.3.2. Server

E.181.3.2.1. Settings

• Convert many `postgresql.conf` settings to enumerated values so that `pg_settings` can display the valid values (Magnus)

• Add `cursor_tuple_fraction` parameter to control the fraction of a cursor’s rows that the planner assumes will be fetched (Robert Hell)

• Allow underscores in the names of custom variable classes in `postgresql.conf` (Tom)

E.181.3.2.2. Authentication and security

• Remove support for the (insecure) `crypt` authentication method (Magnus)

  This effectively obsoletes pre-PostgreSQL 7.2 client libraries, as there is no longer any non-plaintext password method that they can use.

• Support regular expressions in `pg_ident.conf` (Magnus)

• Allow Kerberos/GSSAPI parameters to be changed without restarting the postmaster (Magnus)

• Support SSL certificate chains in server certificate file (Andrew Gierth)

  Including the full certificate chain makes the client able to verify the certificate without having all intermediate CA certificates present in the local store, which is often the case for commercial CAs.

• Report appropriate error message for combination of MD5 authentication and `db_user_name-space` enabled (Bruce)

E.181.3.2.3. `pg_hba.conf`

• Change all authentication options to use `name=value` syntax (Magnus)

  This makes incompatible changes to the `ldap`, `pam` and `ident` authentication methods. All `pg_hba.conf` entries with these methods need to be rewritten using the new format.

• Remove the `ident sameuser` option, instead making that behavior the default if no usermap is specified (Magnus)

• Allow a usermap parameter for all external authentication methods (Magnus)

  Previously a usermap was only supported for `ident` authentication.
• Add `clientcert` option to control requesting of a client certificate (Magnus)

  Previously this was controlled by the presence of a root certificate file in the server's data directory.

• Add `cert` authentication method to allow user authentication via SSL certificates (Magnus)

  Previously SSL certificates could only verify that the client had access to a certificate, not authenticate a user.

• Allow `krb5`, `gssapi` and `sspi` realm and `krb5` host settings to be specified in `pg_hba.conf` (Magnus)

  These override the settings in `postgresql.conf`.

• Add `include_realm` parameter for `krb5`, `gssapi`, and `sspi` methods (Magnus)

  This allows identical usernames from different realms to be authenticated as different database users using usermaps.

• Parse `pg_hba.conf` fully when it is loaded, so that errors are reported immediately (Magnus)

  Previously, most errors in the file wouldn’t be detected until clients tried to connect, so an erroneous file could render the system unusable. With the new behavior, if an error is detected during reload then the bad file is rejected and the postmaster continues to use its old copy.

• Show all parsing errors in `pg_hba.conf` instead of aborting after the first one (Selena Deckelmann)

• Support `ident` authentication over Unix-domain sockets on Solaris (Garick Hamlin)

**E.181.3.2.4. Continuous Archiving**

• Provide an option to `pg_start_backup()` to force its implied checkpoint to finish as quickly as possible (Tom)

  The default behavior avoids excess I/O consumption, but that is pointless if no concurrent query activity is going on.

• Make `pg_stop_backup()` wait for modified WAL files to be archived (Simon)

  This guarantees that the backup is valid at the time `pg_stop_backup()` completes.

• When archiving is enabled, rotate the last WAL segment at shutdown so that all transactions can be archived immediately (Guillaume Smet, Heikki)

• Delay “smart” shutdown while a continuous archiving base backup is in progress (Laurenz Albe)

• Cancel a continuous archiving base backup if “fast” shutdown is requested (Laurenz Albe)

• Allow `recovery.conf` boolean variables to take the same range of string values as `postgresql.conf` boolean variables (Bruce)

**E.181.3.2.5. Monitoring**

• Add `pg_conf_load_time()` to report when the PostgreSQL configuration files were last loaded (George Gensure)

• Add `pg_terminate_backend()` to safely terminate a backend (the SIGTERM signal works also) (Tom, Bruce)

  While it's always been possible to SIGTERM a single backend, this was previously considered unsupported; and testing of the case found some bugs that are now fixed.
Release Notes

- Add ability to track user-defined functions’ call counts and runtimes (Martin Pihlak)

Function statistics appear in a new system view, `pg_stat_user_functions`. Tracking is controlled by the new parameter `track_functions`.

- Allow specification of the maximum query string size in `pg_stat_activity` via new `track_activity_query_size` parameter (Thomas Lee)

- Increase the maximum line length sent to syslog, in hopes of improving performance (Tom)

- Add read-only configuration variables `segment_size`, `wal_block_size`, and `wal_segment_size` (Bernd Helmle)

- When reporting a deadlock, report the text of all queries involved in the deadlock to the server log (Itagaki Takahiro)

- Add `pg_stat_get_activity(pid)` function to return information about a specific process id (Magnus)

- Allow the location of the server’s statistics file to be specified via `stats_temp_directory` (Magnus)

This allows the statistics file to be placed in a RAM-resident directory to reduce I/O requirements. On startup/shutdown, the file is copied to its traditional location (`$PGDATA/global/`) so it is preserved across restarts.

### E.181.3.3. Queries

- Add support for `WINDOW` functions (Hitoshi Harada)

- Add support for `WITH` clauses (CTEs), including `WITH RECURSIVE` (Yoshiyuki Asaba, Tatsuo Ishii, Tom)

- Add `TABLE` command (Peter)

  `TABLE` `tablename` is a SQL standard short-hand for `SELECT * FROM tablename`.

- Allow `AS` to be optional when specifying a `SELECT` (or `RETURNING`) column output label (Hiroshi Saito)

  This works so long as the column label is not any PostgreSQL keyword; otherwise `AS` is still needed.

- Support set-returning functions in `SELECT` result lists even for functions that return their result via a tuplestore (Tom)

  In particular, this means that functions written in PL/pgSQL and other PL languages can now be called this way.

- Support set-returning functions in the output of aggregation and grouping queries (Tom)

- Allow `SELECT FOR UPDATE/SHARE` to work on inheritance trees (Tom)

- Add infrastructure for SQL/MED (Martin Pihlak, Peter)

  There are no remote or external SQL/MED capabilities yet, but this change provides a standardized and future-proof system for managing connection information for modules like `dblink` and `pl-proxy`.

- Invalidate cached plans when referenced schemas, functions, operators, or operator classes are modified (Martin Pihlak, Tom)

  This improves the system's ability to respond to on-the-fly DDL changes.
• Allow comparison of composite types and allow arrays of anonymous composite types (Tom)

This allows constructs such as row(1, 1.1) = any (array[row(7, 7.7), row(1, 1.0)]). This is particularly useful in recursive queries.

• Add support for Unicode string literal and identifier specifications using code points, e.g. \0061t\+000061' (Peter)

• Reject \000 in string literals and COPY data (Tom)

Previously, this was accepted but had the effect of terminating the string contents.

• Improve the parser’s ability to report error locations (Tom)

An error location is now reported for many semantic errors, such as mismatched datatypes, that previously could not be localized.

E.181.3.3.1. TRUNCATE

• Support statement-level ON TRUNCATE triggers (Simon)

• Add RESTART/CONTINUE IDENTITY options for TRUNCATE TABLE (Zoltan Boszormenyi)

The start value of a sequence can be changed by ALTER SEQUENCE START WITH.

• Allow TRUNCATE tab1, tab1 to succeed (Bruce)

• Add a separate TRUNCATE permission (Robert Haas)

E.181.3.3.2. EXPLAIN

• Make EXPLAIN VERBOSE show the output columns of each plan node (Tom)

Previously EXPLAIN VERBOSE output an internal representation of the query plan. (That behavior is now available via debug_print_plan.)

• Make EXPLAIN identify subplans and initplans with individual labels (Tom)

• Make EXPLAIN honor debug_print_plan (Tom)

• Allow EXPLAIN on CREATE TABLE AS (Peter)

E.181.3.3.3. LIMIT/OFFSET

• Allow sub-selects in LIMIT and OFFSET (Tom)

• Add SQL-standard syntax for LIMIT/OFFSET capabilities (Peter)

To wit, OFFSET num (ROW|ROWS) FETCH (FIRST|NEXT) [num] (ROW|ROWS) ONLY.

E.181.3.4. Object Manipulation

• Add support for column-level privileges (Stephen Frost, KaiGai Kohei)

• Refactor multi-object DROP operations to reduce the need for CASCADE (Alex Hunsaker)

For example, if table B has a dependency on table A, the command DROP TABLE A, B no longer requires the CASCADE option.

• Fix various problems with concurrent DROP commands by ensuring that locks are taken before we begin to drop dependencies of an object (Tom)

• Improve reporting of dependencies during DROP commands (Tom)
• Add WITH [NO] DATA clause to CREATE TABLE AS, per the SQL standard (Peter, Tom)

• Add support for user-defined I/O conversion casts (Heikki)

• Allow CREATE AGGREGATE to use an internal transition datatype (Tom)

• Add LIKE clause to CREATE TYPE (Tom)

  This simplifies creation of data types that use the same internal representation as an existing type.

• Allow specification of the type category and “preferred” status for user-defined base types (Tom)

  This allows more control over the coercion behavior of user-defined types.

• Allow CREATE OR REPLACE VIEW to add columns to the end of a view (Robert Haas)

E.181.3.4.1. ALTER

• Add ALTER TYPE RENAME (Petr Jelinek)

• Add ALTER SEQUENCE ... RESTART (with no parameter) to reset a sequence to its initial value (Zoltan Boszormenyi)

• Modify the ALTER TABLE syntax to allow all reasonable combinations for tables, indexes, sequences, and views (Tom)

  This change allows the following new syntaxes:

  • ALTER SEQUENCE OWNER TO

  • ALTER VIEW ALTER COLUMN SET/DROP DEFAULT

  • ALTER VIEW OWNER TO

  • ALTER VIEW SET SCHEMA

  There is no actual new functionality here, but formerly you had to say ALTER TABLE to do these things, which was confusing.

• Add support for the syntax ALTER TABLE ... ALTER COLUMN ... SET DATA TYPE (Peter)

  This is SQL-standard syntax for functionality that was already supported.

• Make ALTER TABLE SET WITHOUT OIDS rewrite the table to physically remove OID values (Tom)

  Also, add ALTER TABLE SET WITH OIDS to rewrite the table to add OIDS.

E.181.3.4.2. Database Manipulation

• Improve reporting of CREATE/DROP/RENAME DATABASE failure when uncommitted prepared transactions are the cause (Tom)

• Make LC_COLLATE and LC_CTYPE into per-database settings (Radek Strnad, Heikki)

  This makes collation similar to encoding, which was always configurable per database.

• Improve checks that the database encoding, collation (LC_COLLATE), and character classes (LC_CTYPE) match (Heikki, Tom)

  Note in particular that a new database's encoding and locale settings can be changed only when copying from template0. This prevents possibly copying data that doesn't match the settings.
• Add ALTER DATABASE SET TABLESPACE to move a database to a new tablespace (Guillaume Lelarge, Bernd Helmle)

E.181.3.5. Utility Operations

• Add a VERBOSE option to the CLUSTER command and clusterdb (Jim Cox)
• Decrease memory requirements for recording pending trigger events (Tom)

E.181.3.5.1. Indexes

• Dramatically improve the speed of building and accessing hash indexes (Tom Raney, Shreya Bhar-gava)

This allows hash indexes to be sometimes faster than btree indexes. However, hash indexes are still not crash-safe.

• Make hash indexes store only the hash code, not the full value of the indexed column (Xiao Meng)

This greatly reduces the size of hash indexes for long indexed values, improving performance.

• Implement fast update option for GIN indexes (Teodor, Oleg)

This option greatly improves update speed at a small penalty in search speed.

• xxx_pattern_ops indexes can now be used for simple equality comparisons, not only for LIKE (Tom)

E.181.3.5.2. Full Text Indexes

• Remove the requirement to use @@ when doing GIN weighted lookups on full text indexes (Tom, Teodor)

The normal @@ text search operator can be used instead.

• Add an optimizer selectivity function for @@ text search operations (Jan Urbanski)
• Allow prefix matching in full text searches (Teodor Sigaev, Oleg Bartunov)
• Support multi-column GIN indexes (Teodor Sigaev)
• Improve support for Nepali language and Devanagari alphabet (Teodor)

E.181.3.5.3. VACUUM

• Track free space in separate per-relation “fork” files (Heikki)

Free space discovered by VACUUM is now recorded in *_fsm files, rather than in a fixed-sized shared memory area. The max_fsm_pages and max_fsm_relations settings have been removed, greatly simplifying administration of free space management.

• Add a visibility map to track pages that do not require vacuuming (Heikki)

This allows VACUUM to avoid scanning all of a table when only a portion of the table needs vacu-uming. The visibility map is stored in per-relation “fork” files.

• Add vacuum_freeze_table_age parameter to control when VACUUM should ignore the vis-ibility map and do a full table scan to freeze tuples (Heikki)

• Track transaction snapshots more carefully (Alvaro)

This improves VACUUM's ability to reclaim space in the presence of long-running transactions.
• Add ability to specify per-relation autovacuum and TOAST parameters in `CREATE TABLE` (Alvaro, Euler Taveira de Oliveira)

Autovacuum options used to be stored in a system table.
• Add `--freeze` option to vacuumdb (Bruce)

E.181.3.6. Data Types

• Add a `CaseSensitive` option for text search synonym dictionaries (Simon)

• Improve the precision of `NUMERIC` division (Tom)

• Add basic arithmetic operators for `int2` with `int8` (Tom)

  This eliminates the need for explicit casting in some situations.

• Allow `UUID` input to accept an optional hyphen after every fourth digit (Robert Haas)

• Allow `on`/`off` as input for the boolean data type (Itagaki Takahiro)

• Allow spaces around `NaN` in the input string for type `numeric` (Sam Mason)

E.181.3.6.1. Temporal Data Types

• Reject year 0 BC and years 000 and 0000 (Tom)

  Previously these were interpreted as 1 BC. (Note: years 0 and 00 are still assumed to be the year 2000.)

• Include `SGT` (Singapore time) in the default list of known time zone abbreviations (Tom)

• Support `infinity` and `-infinity` as values of type `date` (Tom)

• Make parsing of `interval` literals more standard-compliant (Tom, Ron Mayer)

  For example, `INTERVAL '1' YEAR` now does what it's supposed to.

• Allow `interval` fractional-seconds precision to be specified after the `second` keyword, for SQL standard compliance (Tom)

  Formerly the precision had to be specified after the keyword `interval`. (For backwards compatibility, this syntax is still supported, though deprecated.) Data type definitions will now be output using the standard format.

• Support the ISO 8601 `interval` syntax (Ron Mayer, Kevin Grittner)

  For example, `INTERVAL 'P1Y2M3DT4H5M6.7S'` is now supported.

• Add `IntervalStyle` parameter which controls how `interval` values are output (Ron Mayer)

  Valid values are: `postgres`, `postgres_verbose`, `sql_standard`, `iso_8601`. This setting also controls the handling of negative `interval` input when only some fields have positive/negative designations.

• Improve consistency of handling of fractional seconds in `timestamp` and `interval` output (Ron Mayer)

E.181.3.6.2. Arrays

• Improve the handling of casts applied to `ARRAY[]` constructs, such as `ARRAY[...]]::integer[]` (Brendan Jurd)
Formerly PostgreSQL attempted to determine a data type for the \texttt{ARRAY[]} construct without reference to the ensuing cast. This could fail unnecessarily in many cases, in particular when the \texttt{ARRAY[]} construct was empty or contained only ambiguous entries such as \texttt{NULL}. Now the cast is consulted to determine the type that the array elements must be.

- Make SQL-syntax \texttt{ARRAY} dimensions optional to match the SQL standard (Peter)
- Add \texttt{array_ndims()} to return the number of dimensions of an array (Robert Haas)
- Add \texttt{array_length()} to return the length of an array for a specified dimension (Jim Nasby, Robert Haas, Peter Eisentraut)
- Add aggregate function \texttt{array_agg()}, which returns all aggregated values as a single array (Robert Haas, Jeff Davis, Peter)
- Add \texttt{unnest()}, which converts an array to individual row values (Tom)
  This is the opposite of \texttt{array_agg()}.
- Add \texttt{array_fill()} to create arrays initialized with a value (Pavel Stehule)
- Add \texttt{generate_subscripts()} to simplify generating the range of an array's subscripts (Pavel Stehule)

**E.181.3.6.3. Wide-Value Storage (TOAST)**

- Consider TOAST compression on values as short as 32 bytes (previously 256 bytes) (Greg Stark)
- Require 25\% minimum space savings before using TOAST compression (previously 20\% for small values and any-savings-at-all for large values) (Greg)
- Improve TOAST heuristics for rows that have a mix of large and small toastable fields, so that we prefer to push large values out of line and don’t compress small values unnecessarily (Greg, Tom)

**E.181.3.7. Functions**

- Document that \texttt{setseed()} allows values from $-1$ to $1$ (not just $0$ to $1$), and enforce the valid range (Kris Jurka)
- Add server-side function \texttt{lo_import(filename, oid)} (Tatsuo)
- Add \texttt{quote_nullable()}, which behaves like \texttt{quote_literal()} but returns the string \texttt{NULL} for a null argument (Brendan Jurd)
- Improve full text search \texttt{headline()} function to allow extracting several fragments of text (Sushant Sinha)
- Add \texttt{suppress_redundant_updates_trigger()} trigger function to avoid overhead for non-data-changing updates (Andrew)
- Add \texttt{div(numeric, numeric)} to perform numeric division without rounding (Tom)
- Add \texttt{timestamp} and \texttt{timestamptz} versions of \texttt{generate_series()} (Hitoshi Harada)

**E.181.3.7.1. Object Information Functions**

- Implement \texttt{current_query()} for use by functions that need to know the currently running query (Tomas Doran)
- Add \texttt{pg_get_keywords()} to return a list of the parser keywords (Dave Page)
- Add \texttt{pg_get_functiondef()} to see a function's definition (Abhijit Menon-Sen)
• Allow the second argument of `pg_get_expr()` to be zero when deparsing an expression that does not contain variables (Tom)

• Modify `pg_relation_size()` to use `regclass` (Heikki)

  `pg_relation_size(data_type_name)` no longer works.

• Add `boot_val` and `reset_val` columns to `pg_settings` output (Greg Smith)

• Add source file name and line number columns to `pg_settings` output for variables set in a configuration file (Magnus, Alvaro)

  For security reasons, these columns are only visible to superusers.

• Add support for `CURRENT_CATALOG`, `CURRENT_SCHEMA`, `SET CATALOG`, `SET SCHEMA` (Peter)

  These provide SQL-standard syntax for existing features.

• Add `pg_typeof()` which returns the data type of any value (Brendan Jurd)

• Make `version()` return information about whether the server is a 32- or 64-bit binary (Bruce)

• Fix the behavior of information schema columns `is_insertable_into` and `is_updatable` to be consistent (Peter)

• Improve the behavior of information schema `datetime_precision` columns (Peter)

  These columns now show zero for date columns, and 6 (the default precision) for time, timestamp, and interval without a declared precision, rather than showing null as formerly.

• Convert remaining builtin set-returning functions to use `OUT` parameters (Jaime Casanova)

  This makes it possible to call these functions without specifying a column list: `pg_show_all_settings()`, `pg_lock_status()`, `pg_prepared_xact()`, `pg_prepared_statement()`, `pg_cursor()`.

• Make `pg_*_is_visible()` and `has_*_privilege()` functions return NULL for invalid OIDs, rather than reporting an error (Tom)

• Extend `has_*_privilege()` functions to allow inquiring about the OR of multiple privileges in one call (Stephen Frost, Tom)

• Add `has_column_privilege()` and `has_any_column_privilege()` functions (Stephen Frost, Tom)

### E.181.3.7.2. Function Creation

• Support variadic functions (functions with a variable number of arguments) (Pavel Stehule)

  Only trailing arguments can be optional, and they all must be of the same data type.

• Support default values for function arguments (Pavel Stehule)

• Add `CREATE FUNCTION ... RETURNS TABLE` clause (Pavel Stehule)

• Allow SQL-language functions to return the output of an `INSERT/UPDATE/DELETE RETURNING` clause (Tom)

### E.181.3.7.3. PL/pgSQL Server-Side Language

• Support `EXECUTE USING` for easier insertion of data values into a dynamic query string (Pavel Stehule)
• Allow looping over the results of a cursor using a FOR loop (Pavel Stehule)

• Support RETURN QUERY EXECUTE (Pavel Stehule)

• Improve the RAISE command (Pavel Stehule)
  • Support DETAIL and HINT fields
  • Support specification of the SQLSTATE error code
  • Support an exception name parameter
  • Allow RAISE without parameters in an exception block to re-throw the current error

• Allow specification of SQLSTATE codes in EXCEPTION lists (Pavel Stehule)
  This is useful for handling custom SQLSTATE codes.

• Support the CASE statement (Pavel Stehule)

• Make RETURN QUERY set the special FOUND and GET DIAGNOSTICS ROW_COUNT variables (Pavel Stehule)

• Make FETCH and MOVE set the GET DIAGNOSTICS ROW_COUNT variable (Andrew Gierth)

• Make EXIT without a label always exit the innermost loop (Tom)
  Formerly, if there were a BEGIN block more closely nested than any loop, it would exit that block instead. The new behavior matches Oracle(TM) and is also what was previously stated by our own documentation.

• Make processing of string literals and nested block comments match the main SQL parser’s processing (Tom)
  In particular, the format string in RAISE now works the same as any other string literal, including being subject to standard_conforming_strings. This change also fixes other cases in which valid commands would fail when standard_conforming_strings is on.

• Avoid memory leakage when the same function is called at varying exception-block nesting depths (Tom)

E.181.3.8. Client Applications

• Fix pg_ctl restart to preserve command-line arguments (Bruce)

• Add -w/--no-password option that prevents password prompting in all utilities that have a -W/--password option (Peter)

• Remove -q (quiet) option of createdb, createuser, dropdb, dropuser (Peter)
  These options have had no effect since PostgreSQL 8.3.

E.181.3.8.1. psql

• Remove verbose startup banner; now just suggest help (Joshua Drake)

• Make help show common backlash commands (Greg Sabino Mullane)

• Add \pset format wrapped mode to wrap output to the screen width, or file/pipe output too if \pset columns is set (Bryce Nesbitt)

• Allow all supported spellings of boolean values in \pset, rather than just on and off (Bruce)
Formerly, any string other than “off” was silently taken to mean \texttt{true}. \texttt{psql} will now complain about unrecognized spellings (but still take them as \texttt{true}).

- Use the pager for wide output (Bruce)
- Require a space between a one-letter backslash command and its first argument (Bernd Helmle)
  This removes a historical source of ambiguity.
- Improve tab completion support for schema-qualified and quoted identifiers (Greg Sabino Mullane)
- Add optional \texttt{on/off} argument for \texttt{\textbackslash timing} (David Fetter)
- Display access control rights on multiple lines (Brendan Jurd, Andreas Scherbaum)
- Make \texttt{\textbackslash l} show database access privileges (Andrew Gilligan)
- Make \texttt{\textbackslash l+} show database sizes, if permissions allow (Andrew Gilligan)
- Add the \texttt{\textbackslash ef} command to edit function definitions (Abhijit Menon-Sen)

\textbf{E.181.3.8.2. psql \textbackslash d* commands}

- Make \texttt{\textbackslash d*} commands that do not have a pattern argument show system objects only if the \texttt{S} modifier is specified (Greg Sabino Mullane, Bruce)

  The former behavior was inconsistent across different variants of \texttt{\textbackslash d}, and in most cases it provided no easy way to see just user objects.

- Improve \texttt{\textbackslash d*} commands to work with older PostgreSQL server versions (back to 7.4), not only the current server version (Guillaume Lelarge)

- Make \texttt{\textbackslash d} show foreign-key constraints that reference the selected table (Kenneth D'Souza)

- Make \texttt{\textbackslash d} on a sequence show its column values (Euler Taveira de Oliveira)

- Add column storage type and other relation options to the \texttt{\textbackslash d+} display (Gregory Stark, Euler Taveira de Oliveira)

- Show relation size in \texttt{\textbackslash d t+} output (Dickson S. Guedes)

- Show the possible values of \texttt{enum} types in \texttt{\textbackslash d T+} (David Fetter)

- Allow \texttt{\textbackslash d C} to accept a wildcard pattern, which matches either datatype involved in the cast (Tom)

- Add a function type column to \texttt{\textbackslash d f}'s output, and add options to list only selected types of functions (David Fetter)

- Make \texttt{\textbackslash d f} not hide functions that take or return type \texttt{cstring} (Tom)

  Previously, such functions were hidden because most of them are datatype I/O functions, which were deemed uninteresting. The new policy about hiding system functions by default makes this \texttt{wrt} unnecessary.

\textbf{E.181.3.8.3. pg_dump}

- Add a \texttt{--no-tablespaces} option to \texttt{pg_dump/pg_dumpall/pg_restore} so that dumps can be restored to clusters that have non-matching tablespace layouts (Gavin Roy)

- Remove \texttt{-d} and \texttt{-D} options from \texttt{pg_dump} and \texttt{pg_dumpall} (Tom)
These options were too frequently confused with the option to select a database name in other PostgreSQL client applications. The functionality is still available, but you must now spell out the long option name `--inserts` or `--column-inserts`.

- Remove `-i|--ignore-version` option from `pg_dump` and `pg_dumpall` (Tom)
  
  Use of this option does not throw an error, but it has no effect. This option was removed because the version checks are necessary for safety.

- Disable `statement_timeout` during dump and restore (Joshua Drake)

- Add `pg_dump/pg_dumpall` option `--lock-wait-timeout` (David Gould)
  
  This allows dumps to fail if unable to acquire a shared lock within the specified amount of time.

- Reorder `pg_dump --data-only` output to dump tables referenced by foreign keys before the referencing tables (Tom)
  
  This allows data loads when foreign keys are already present. If circular references make a safe ordering impossible, a `NOTICE` is issued.

- Allow `pg_dump`, `pg_dumpall`, and `pg_restore` to use a specified role (Benedek László)

- Allow `pg_restore` to use multiple concurrent connections to do the restore (Andrew)
  
  The number of concurrent connections is controlled by the option `--jobs`. This is supported only for custom-format archives.

**E.181.3.9. Programming Tools**

**E.181.3.9.1. libpq**

- Allow the OID to be specified when importing a large object, via new function `lo_import_with_oid()` (Tatsuo)

- Add “events” support (Andrew Chernow, Merlin Moncure)
  
  This adds the ability to register callbacks to manage private data associated with `PGconn` and `PGresult` objects.

- Improve error handling to allow the return of multiple error messages as multi-line error reports (Magnus)

- Make `PQexecParams()` and related functions return `PGRES_EMPTY_QUERY` for an empty query (Tom)
  
  They previously returned `PGRES_COMMAND_OK`.

- Document how to avoid the overhead of `WSACleanup()` on Windows (Andrew Chernow)

- Do not rely on Kerberos tickets to determine the default database username (Magnus)
  
  Previously, a Kerberos-capable build of `libpq` would use the principal name from any available Kerberos ticket as default database username, even if the connection wasn't using Kerberos authentication. This was deemed inconsistent and confusing. The default username is now determined the same way with or without Kerberos. Note however that the database username must still match the ticket when Kerberos authentication is used.

**E.181.3.9.2. libpq SSL (Secure Sockets Layer) support**

- Fix certificate validation for SSL connections (Magnus)
libpq now supports verifying both the certificate and the name of the server when making SSL
connections. If a root certificate is not available to use for verification, SSL connections will fail.
The `sslmode` parameter is used to enable certificate verification and set the level of checking.
The default is still not to do any verification, allowing connections to SSL-enabled servers without
requiring a root certificate on the client.

- Support wildcard server certificates (Magnus)
  
  If a certificate CN starts with `*`, it will be treated as a wildcard when matching the hostname,
  allowing the use of the same certificate for multiple servers.

- Allow the file locations for client certificates to be specified (Mark Woodward, Alvaro, Magnus)

- Add a `PQinitOpenSSL` function to allow greater control over OpenSSL/libcrypto initialization
  (Andrew Chernow)

- Make libpq unregister its OpenSSL callbacks when no database connections remain open (Bruce,
  Magnus, Russell Smith)
  
  This is required for applications that unload the libpq library, otherwise invalid OpenSSL callbacks
  will remain.

E.181.3.9.3. ecpg

- Add localization support for messages (Euler Taveira de Oliveira)

- `ecpg` parser is now automatically generated from the server parser (Michael)
  
  Previously the `ecpg` parser was hand-maintained.

E.181.3.9.4. Server Programming Interface (SPI)

- Add support for single-use plans with out-of-line parameters (Tom)

- Add new `SPI_OK_REWRITTEN` return code for `SPI_execute()` (Heikki)
  
  This is used when a command is rewritten to another type of command.

- Remove unnecessary inclusions from `executor/spi.h` (Tom)
  
  SPI-using modules might need to add some `#include` lines if they were depending on `spi.h`
  to include things for them.

E.181.3.10. Build Options

- Update build system to use Autoconf 2.61 (Peter)

- Require GNU bison for source code builds (Peter)
  
  This has effectively been required for several years, but now there is no infrastructure claiming to
  support other parser tools.

- Add `pg_config --htmldir` option (Peter)

- Pass `float4` by value inside the server (Zoltan Boszormenyi)
  
  Add configure option `--disable-float4-byval` to use the old behavior. External C functions
  that use old-style (version 0) call convention and pass or return `float4` values will be broken by
  this change, so you may need the configure option if you have such functions and don’t want to
  update them.
• Pass float8, int8, and related datatypes by value inside the server on 64-bit platforms (Zoltan Boszormenyi)

Add configure option --disable-float8-byval to use the old behavior. As above, this change might break old-style external C functions.

• Add configure options --with-segsize, --with-blocksize, --with-wal-block-size, --with-wal-segsize (Zdenek Kotala, Tom)

This simplifies build-time control over several constants that previously could only be changed by editing pg_config_manual.h.

• Allow threaded builds on Solaris 2.5 (Bruce)

• Use the system's getopt_long() on Solaris (Zdenek Kotala, Tom)

This makes option processing more consistent with what Solaris users expect.

• Add support for the Sun Studio compiler on Linux (Julius Strofek)

• Append the major version number to the backend gettext domain, and the soname major version number to libraries' gettext domain (Peter)

This simplifies parallel installations of multiple versions.

• Add support for code coverage testing with gcov (Michelle Caisse)

• Allow out-of-tree builds on Mingw and Cygwin (Richard Evans)

• Fix the use of Mingw as a cross-compiling source platform (Peter)

E.181.3.11. Source Code

• Support 64-bit time zone data files (Heikki)

This adds support for daylight saving time (DST) calculations beyond the year 2038.

• Deprecate use of platform's time_t data type (Tom)

Some platforms have migrated to 64-bit time_t, some have not, and Windows can't make up its mind what it's doing. Define pg_time_t to have the same meaning as time_t, but always be 64 bits (unless the platform has no 64-bit integer type), and use that type in all module APIs and on-disk data formats.

• Fix bug in handling of the time zone database when cross-compiling (Richard Evans)

• Link backend object files in one step, rather than in stages (Peter)

• Improve gettext support to allow better translation of plurals (Peter)

• Add message translation support to the PL languages (Alvaro, Peter)

• Add more DTrace probes (Robert Lor)

• Enable DTrace support on macOS Leopard and other non-Solaris platforms (Robert Lor)

• Simplify and standardize conversions between C strings and text datums, by providing common functions for the purpose (Brendan Jurd, Tom)

• Clean up the include/catalog/ header files so that frontend programs can include them without including postgres.h (Zdenek Kotala)

• Make name char-aligned, and suppress zero-padding of name entries in indexes (Tom)
• Recover better if dynamically-loaded code executes exit() (Tom)

• Add a hook to let plug-ins monitor the executor (Itagaki Takahiro)

• Add a hook to allow the planner’s statistics lookup behavior to be overridden (Simon Riggs)

• Add shmem_startup_hook() for custom shared memory requirements (Tom)

• Replace the index access method amgetmulti entry point with amgetbitmap, and extend the API for amgettuple to support run-time determination of operator lossiness (Heikki, Tom, Teodor)

The API for GIN and GiST opclass consistent functions has been extended as well.

• Add support for partial-match searches in GIN indexes (Teodor Sigaev, Oleg Bartunov)

• Replace pg_class column reltriggers with boolean relhastriggers (Simon)
  Also remove unused pg_class columns relukeys, relfkeys, and relrefs.

• Add a relistemp column to pg_class to ease identification of temporary tables (Tom)

• Move platform FAQs into the main documentation (Peter)

• Prevent parser input files from being built with any conflicts (Peter)

• Add support for the KOI8U (Ukrainian) encoding (Peter)

• Add Japanese message translations (Japan PostgreSQL Users Group)
  This used to be maintained as a separate project.

• Fix problem when setting LC_MESSAGES on MSVC-built systems (Hiroshi Inoue, Hiroshi Saito, Magnus)

E.181.3.12. Contrib

• Add contrib/auto_explain to automatically run EXPLAIN on queries exceeding a specified duration (Itagaki Takahiro, Tom)

• Add contrib/btree_gin to allow GIN indexes to handle more datatypes (Oleg, Teodor)

• Add contrib/citext to provide a case-insensitive, multibyte-aware text data type (David Wheeler)

• Add contrib/pg_stat_statements for server-wide tracking of statement execution statistics (Itagaki Takahiro)

• Add duration and query mode options to contrib/pgbench (Itagaki Takahiro)

• Make contrib/pgbench use table names pgbench_accounts, pgbench_branches, pgbench_history, and pgbench_tellers, rather than just accounts, branches, history, and tellers (Tom)
  This is to reduce the risk of accidentally destroying real data by running pgbench.

• Fix contrib/pgstattuple to handle tables and indexes with over 2 billion pages (Tatsuhito Kasahara)

• In contrib/fuzzystrmatch, add a version of the Levenshtein string-distance function that allows the user to specify the costs of insertion, deletion, and substitution (Volkan Yazici)

• Make contrib/ltree support multibyte encodings (laser)
• Enable contrib/dblink to use connection information stored in the SQL/MED catalogs (Joe Conway)

• Improve contrib/dblink's reporting of errors from the remote server (Joe Conway)

• Make contrib/dblink set client_encoding to match the local database's encoding (Joe Conway)

  This prevents encoding problems when communicating with a remote database that uses a different encoding.

• Make sure contrib/dblink uses a password supplied by the user, and not accidentally taken from the server's .pgpass file (Joe Conway)

  This is a minor security enhancement.

• Add fsm_page_contents() to contrib/pageinspect (Heikki)

• Modify get_raw_page() to support free space map (*_fsm) files. Also update contrib/pg_freespacemap.

• Add support for multibyte encodings to contrib/pg_trgm (Teodor)

• Rewrite contrib/intagg to use new functions array_agg() and unnest() (Tom)

• Make contrib/pg_standby recover all available WAL before failover (Fujii Masao, Simon, Heikki)

  To make this work safely, you now need to set the new recovery_end_command option in recovery.conf to clean up the trigger file after failover. pg_standby will no longer remove the trigger file itself.

• contrib/pg_standby's -l option is now a no-op, because it is unsafe to use a symlink (Simon)

E.182. Release 8.3.23

Release date: 2013-02-07

This release contains a variety of fixes from 8.3.22. For information about new features in the 8.3 major release, see Section E.205.

This is expected to be the last PostgreSQL release in the 8.3.X series. Users are encouraged to update to a newer release branch soon.

E.182.1. Migration to Version 8.3.23

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

E.182.2. Changes

• Prevent execution of enum_recv from SQL (Tom Lane)

  The function was misdeclared, allowing a simple SQL command to crash the server. In principle an attacker might be able to use it to examine the contents of server memory. Our thanks to Sumit Soni (via Secunia SVCRP) for reporting this issue. (CVE-2013-0255)

• Fix SQL grammar to allow subscripting or field selection from a sub-SELECT result (Tom Lane)
• Protect against race conditions when scanning `pg_tablespace` (Stephen Frost, Tom Lane)

CREATE DATABASE and DROP DATABASE could misbehave if there were concurrent updates of `pg_tablespace` entries.

• Prevent `DROP OWNED` from trying to drop whole databases or tablespaces (Álvaro Herrera)

For safety, ownership of these objects must be reassigned, not dropped.

• Prevent misbehavior when a `RowExpr` or `XmlExpr` is parse-analyzed twice (Andres Freund, Tom Lane)

This mistake could be user-visible in contexts such as CREATE TABLE LIKE INCLUDING INDEXES.

• Improve defenses against integer overflow in hashtable sizing calculations (Jeff Davis)

• Ensure that non-ASCII prompt strings are translated to the correct code page on Windows (Alexander Law, Noah Misch)

This bug affected psql and some other client programs.

• Fix possible crash in psql's `\?` command when not connected to a database (Meng Qingzhong)

• Fix one-byte buffer overrun in libpq's `PQprintTuples` (Xi Wang)

This ancient function is not used anywhere by PostgreSQL itself, but it might still be used by some client code.

• Rearrange configure's tests for supplied functions so it is not fooled by bogus exports from libedit/libreadline (Christoph Berg)

• Ensure Windows build number increases over time (Magnus Hagander)

• Make `pgxs` build executables with the right `.exe` suffix when cross-compiling for Windows (Zoltan Boszormenyi)

• Add new timezone abbreviation `FET` (Tom Lane)

This is now used in some eastern-European time zones.

E.183. Release 8.3.22

Release date: 2012-12-06

This release contains a variety of fixes from 8.3.21. For information about new features in the 8.3 major release, see Section E.205.

The PostgreSQL community will stop releasing updates for the 8.3.X release series in February 2013. Users are encouraged to update to a newer release branch soon.

E.183.1. Migration to Version 8.3.22

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

E.183.2. Changes

• Fix multiple bugs associated with `CREATE INDEX CONCURRENTLY` (Andres Freund, Tom Lane)
Fix CREATE INDEX CONCURRENTLY to use in-place updates when changing the state of an index's pg_index row. This prevents race conditions that could cause concurrent sessions to miss updating the target index, thus resulting in corrupt concurrently-created indexes.

Also, fix various other operations to ensure that they ignore invalid indexes resulting from a failed CREATE INDEX CONCURRENTLY command. The most important of these is VACUUM, because an auto-vacuum could easily be launched on the table before corrective action can be taken to fix or remove the invalid index.

- Avoid corruption of internal hash tables when out of memory (Hitoshi Harada)
- Fix planning of non-strict equivalence clauses above outer joins (Tom Lane)
  The planner could derive incorrect constraints from a clause equating a non-strict construct to something else, for example WHERE COALESCE(foo, 0) = 0 when foo is coming from the nullable side of an outer join.
- Improve planner's ability to prove exclusion constraints from equivalence classes (Tom Lane)
- Fix partial-row matching in hashed subplans to handle cross-type cases correctly (Tom Lane)
  This affects multicolumn NOT IN subplans, such as WHERE (a, b) NOT IN (SELECT x, y FROM ...) when for instance b and y are int4 and int8 respectively. This mistake led to wrong answers or crashes depending on the specific datatypes involved.
- Acquire buffer lock when re-fetching the old tuple for an AFTER ROW UPDATE/DELETE trigger (Andres Freund)
  In very unusual circumstances, this oversight could result in passing incorrect data to the precheck logic for a foreign-key enforcement trigger. That could result in a crash, or in an incorrect decision about whether to fire the trigger.
- Fix REASSIGN OWNED to handle grants on tablespaces (Álvaro Herrera)
- Ignore incorrect pg_attribute entries for system columns for views (Tom Lane)
  Views do not have any system columns. However, we forgot to remove such entries when converting a table to a view. That's fixed properly for 9.3 and later, but in previous branches we need to defend against existing mis-converted views.
- Fix rule printing to dump INSERT INTO table DEFAULT VALUES correctly (Tom Lane)
- Guard against stack overflow when there are too many UNION/INTERSECT/EXCEPT clauses in a query (Tom Lane)
- Prevent platform-dependent failures when dividing the minimum possible integer value by -1 (Xi Wang, Tom Lane)
- Fix possible access past end of string in date parsing (Hitoshi Harada)
- Produce an understandable error message if the length of the path name for a Unix-domain socket exceeds the platform-specific limit (Tom Lane, Andrew Dunstan)
  Formerly, this would result in something quite unhelpful, such as "Non-recoverable failure in name resolution".
- Fix memory leaks when sending composite column values to the client (Tom Lane)
- Make pg_ctl more robust about reading the postmaster.pid file (Heikki Linnakangas)
  Fix race conditions and possible file descriptor leakage.
• Fix possible crash in psql if incorrectly-encoded data is presented and the `client_encoding` setting is a client-only encoding, such as SJIS (Jiang Guiqing)

• Fix bugs in the `restore.sql` script emitted by `pg_dump` in `tar` output format (Tom Lane)

  The script would fail outright on tables whose names include upper-case characters. Also, make the script capable of restoring data in `--inserts` mode as well as the regular COPY mode.

• Fix `pg_restore` to accept POSIX-conformant `tar` files (Brian Weaver, Tom Lane)

  The original coding of `pg_dump`'s `tar` output mode produced files that are not fully conformant with the POSIX standard. This has been corrected for version 9.3. This patch updates previous branches so that they will accept both the incorrect and the corrected formats, in hopes of avoiding compatibility problems when 9.3 comes out.

• Fix `pg_resetxlog` to locate `postmaster.pid` correctly when given a relative path to the data directory (Tom Lane)

  This mistake could lead to `pg_resetxlog` not noticing that there is an active postmaster using the data directory.

• Fix `libpq`'s `lo_import()` and `lo_export()` functions to report file I/O errors properly (Tom Lane)

• Fix ecpg's processing of nested structure pointer variables (Muhammad Usama)

• Make `contrib/pageinspect`'s btree page inspection functions take buffer locks while examining pages (Tom Lane)

• Fix `pgxs` support for building loadable modules on AIX (Tom Lane)

  Building modules outside the original source tree didn't work on AIX.

• Update time zone data files to `tzdata` release 2012j for DST law changes in Cuba, Israel, Jordan, Libya, Palestine, Western Samoa, and portions of Brazil.

### E.184. Release 8.3.21

**Release date:** 2012-09-24

This release contains a variety of fixes from 8.3.20. For information about new features in the 8.3 major release, see Section E.205.

The PostgreSQL community will stop releasing updates for the 8.3.X release series in February 2013. Users are encouraged to update to a newer release branch soon.

#### E.184.1. Migration to Version 8.3.21

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

#### E.184.2. Changes

• Improve page-splitting decisions in GiST indexes (Alexander Korotkov, Robert Haas, Tom Lane)

  Multi-column GiST indexes might suffer unexpected bloat due to this error.

• Fix cascading privilege revoke to stop if privileges are still held (Tom Lane)
If we revoke a grant option from some role \( X \), but \( X \) still holds that option via a grant from someone else, we should not recursively revoke the corresponding privilege from role(s) \( Y \) that \( X \) had granted it to.

- Fix handling of SIGFPE when PL/Perl is in use (Andres Freund)
  
  Perl resets the process's SIGFPE handler to SIG_IGN, which could result in crashes later on. Restore the normal Postgres signal handler after initializing PL/Perl.

- Prevent PL/Perl from crashing if a recursive PL/Perl function is redefined while being executed (Tom Lane)

- Work around possible misoptimization in PL/Perl (Tom Lane)
  
  Some Linux distributions contain an incorrect version of pthread.h that results in incorrect compiled code in PL/Perl, leading to crashes if a PL/Perl function calls another one that throws an error.

- Update time zone data files to tzdata release 2012f for DST law changes in Fiji

## E.185. Release 8.3.20

**Release date:** 2012-08-17

This release contains a variety of fixes from 8.3.19. For information about new features in the 8.3 major release, see Section E.205.

The PostgreSQL community will stop releasing updates for the 8.3.X release series in February 2013. Users are encouraged to update to a newer release branch soon.

### E.185.1. Migration to Version 8.3.20

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

### E.185.2. Changes

- Prevent access to external files/URLs via XML entity references (Noah Misch, Tom Lane)

  `xml_parse()` would attempt to fetch external files or URLs as needed to resolve DTD and entity references in an XML value, thus allowing unprivileged database users to attempt to fetch data with the privileges of the database server. While the external data wouldn't get returned directly to the user, portions of it could be exposed in error messages if the data didn’t parse as valid XML; and in any case the mere ability to check existence of a file might be useful to an attacker. (CVE-2012-3489)

- Prevent access to external files/URLs via contrib/xml2's `xslt_process()` (Peter Eisen- traut)

  libxml offers the ability to read and write both files and URLs through stylesheet commands, thus allowing unprivileged database users to both read and write data with the privileges of the database server. Disable that through proper use of libxml's security options. (CVE-2012-3488)

  Also, remove `xslt_process()`'s ability to fetch documents and stylesheet from external files/URLs. While this was a documented “feature”, it was long regarded as a bad idea. The fix for CVE-2012-3489 broke that capability, and rather than expend effort on trying to fix it, we're just going to summarily remove it.

- Prevent too-early recycling of btree index pages (Noah Misch)
When we allowed read-only transactions to skip assigning XIDs, we introduced the possibility that a deleted btree page could be recycled while a read-only transaction was still in flight to it. This would result in incorrect index search results. The probability of such an error occurring in the field seems very low because of the timing requirements, but nonetheless it should be fixed.

• Fix crash-safety bug with newly-created-or-reset sequences (Tom Lane)

If ALTER SEQUENCE was executed on a freshly created or reset sequence, and then precisely one nextval() call was made on it, and then the server crashed, WAL replay would restore the sequence to a state in which it appeared that no nextval() had been done, thus allowing the first sequence value to be returned again by the next nextval() call. In particular this could manifest for serial columns, since creation of a serial column's sequence includes an ALTER SEQUENCE OWNED BY step.

• Ensure the backup_label file is fsync'd after pg_start_backup() (Dave Kerr)

• Back-patch 9.1 improvement to compress the fsync request queue (Robert Haas)

This improves performance during checkpoints. The 9.1 change has now seen enough field testing to seem safe to back-patch.

• Only allow autovacuum to be auto-canceled by a directly blocked process (Tom Lane)

The original coding could allow inconsistent behavior in some cases; in particular, an autovacuum could get canceled after less than deadlock_timeout grace period.

• Improve logging of autovacuum cancels (Robert Haas)

• Fix log collector so that log_truncate_on_rotation works during the very first log rotation after server start (Tom Lane)

• Ensure that a whole-row reference to a subquery doesn't include any extra GROUP BY or ORDER BY columns (Tom Lane)

• Disallow copying whole-row references in CHECK constraints and index definitions during CREATE TABLE (Tom Lane)

This situation can arise in CREATE TABLE with LIKE or INHERITS. The copied whole-row variable was incorrectly labeled with the row type of the original table not the new one. Rejecting the case seems reasonable for LIKE, since the row types might well diverge later. For INHERITS we should ideally allow it, with an implicit coercion to the parent table's row type; but that will require more work than seems safe to back-patch.

• Fix memory leak in ARRAY(SELECT ...) subqueries (Heikki Linnakangas, Tom Lane)

• Fix extraction of common prefixes from regular expressions (Tom Lane)

The code could get confused by quantified parenthesized subexpressions, such as ^(foo)?bar. This would lead to incorrect index optimization of searches for such patterns.

• Report errors properly in contrib/xml2's xslt_process() (Tom Lane)

• Update time zone data files to tzdata release 2012e for DST law changes in Morocco and Tokelau

E.186. Release 8.3.19

Release date: 2012-06-04

This release contains a variety of fixes from 8.3.18. For information about new features in the 8.3 major release, see Section E.205.
E.186.1. Migration to Version 8.3.19

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

E.186.2. Changes

- Fix incorrect password transformation in contrib/pgcrypto's DES crypt() function (Solar Designer)

  If a password string contained the byte value 0x80, the remainder of the password was ignored, causing the password to be much weaker than it appeared. With this fix, the rest of the string is properly included in the DES hash. Any stored password values that are affected by this bug will thus no longer match, so the stored values may need to be updated. (CVE-2012-2143)

- Ignore SECURITY DEFINER and SET attributes for a procedural language's call handler (Tom Lane)

  Applying such attributes to a call handler could crash the server. (CVE-2012-2655)

- Allow numeric timezone offsets in timestamp input to be up to 16 hours away from UTC (Tom Lane)

  Some historical time zones have offsets larger than 15 hours, the previous limit. This could result in dumped data values being rejected during reload.

- Fix timestamp conversion to cope when the given time is exactly the last DST transition time for the current timezone (Tom Lane)

  This oversight has been there a long time, but was not noticed previously because most DST-using zones are presumed to have an indefinite sequence of future DST transitions.

- Fix text to name and char to name casts to perform string truncation correctly in multibyte encodings (Karl Schnaitter)

- Fix memory copying bug in to_tsquery() (Heikki Linnakangas)

- Fix slow session startup when pg_attribute is very large (Tom Lane)

  If pg_attribute exceeds one-fourth of shared_buffers, cache rebuilding code that is sometimes needed during session start would trigger the synchronized-scan logic, causing it to take many times longer than normal. The problem was particularly acute if many new sessions were starting at once.

- Ensure sequential scans check for query cancel reasonably often (Merlin Moncure)

  A scan encountering many consecutive pages that contain no live tuples would not respond to interrupts meanwhile.

- Ensure the Windows implementation of PGSemaphoreLock() clears ImmediateInterruptOK before returning (Tom Lane)

  This oversight meant that a query-cancel interrupt received later in the same query could be accepted at an unsafe time, with unpredictable but not good consequences.

- Show whole-row variables safely when printing views or rules (Abbas Butt, Tom Lane)

  Corner cases involving ambiguous names (that is, the name could be either a table or column name of the query) were printed in an ambiguous way, risking that the view or rule would be interpreted differently after dump and reload. Avoid the ambiguous case by attaching a no-op cast.
• Ensure autovacuum worker processes perform stack depth checking properly (Heikki Linnakangas)

Previously, infinite recursion in a function invoked by auto-Analyze could crash worker processes.

• Fix logging collector to not lose log coherency under high load (Andrew Dunstan)

The collector previously could fail to reassemble large messages if it got too busy.

• Fix logging collector to ensure it will restart file rotation after receiving SIGHUP (Tom Lane)

• Fix PL/pgSQL's GET DIAGNOSTICS command when the target is the function's first variable (Tom Lane)

• Fix several performance problems in pg_dump when the database contains many objects (Jeff Janes, Tom Lane)

pg_dump could get very slow if the database contained many schemas, or if many objects are in dependency loops, or if there are many owned sequences.

• Fix contrib/dblink's dblink_exec() to not leak temporary database connections upon error (Tom Lane)

• Update time zone data files to tzdata release 2012c for DST law changes in Antarctica, Armenia, Chile, Cuba, Falkland Islands, Gaza, Haiti, Hebron, Morocco, Syria, and Tokelau Islands; also historical corrections for Canada.

E.187. Release 8.3.18

Release date: 2012-02-27

This release contains a variety of fixes from 8.3.17. For information about new features in the 8.3 major release, see Section E.205.

E.187.1. Migration to Version 8.3.18

A dump/restore is not required for those running 8.3.X.

However, if you are upgrading from a version earlier than 8.3.17, see Section E.188.

E.187.2. Changes

• Require execute permission on the trigger function for CREATE TRIGGER (Robert Haas)

This missing check could allow another user to execute a trigger function with forged input data, by installing it on a table he owns. This is only of significance for trigger functions marked SECURITY DEFINER, since otherwise trigger functions run as the table owner anyway. (CVE-2012-0866)

• Convert newlines to spaces in names written in pg_dump comments (Robert Haas)

pg_dump was incautious about sanitizing object names that are emitted within SQL comments in its output script. A name containing a newline would at least render the script syntactically incorrect. Maliciously crafted object names could present a SQL injection risk when the script is reloaded. (CVE-2012-0868)

• Fix btree index corruption from insertions concurrent with vacuuming (Tom Lane)

An index page split caused by an insertion could sometimes cause a concurrently-running VACUUM to miss removing index entries that it should remove. After the corresponding table rows are
removed, the dangling index entries would cause errors (such as “could not read block N in file ...”) or worse, silently wrong query results after unrelated rows are re-inserted at the now-free table locations. This bug has been present since release 8.2, but occurs so infrequently that it was not diagnosed until now. If you have reason to suspect that it has happened in your database, reindexing the affected index will fix things.

• Allow non-existent values for some settings in ALTER USER/DATABASE SET (Heikki Linnakangas)

Allow default_text_search_config, default_tablespace, and temp_tablespaces to be set to names that are not known. This is because they might be known in another database where the setting is intended to be used, or for the tablespace cases because the tablespace might not be created yet. The same issue was previously recognized for search_path, and these settings now act like that one.

• Track the OID counter correctly during WAL replay, even when it wraps around (Tom Lane)

Previously the OID counter would remain stuck at a high value until the system exited replay mode. The practical consequences of that are usually nil, but there are scenarios wherein a standby server that's been promoted to master might take a long time to advance the OID counter to a reasonable value once values are needed.

• Fix regular expression back-references with * attached (Tom Lane)

Rather than enforcing an exact string match, the code would effectively accept any string that satisfies the pattern sub-expression referenced by the back-reference symbol.

A similar problem still afflicts back-references that are embedded in a larger quantified expression, rather than being the immediate subject of the quantifier. This will be addressed in a future PostgreSQL release.

• Fix recently-introduced memory leak in processing of inet/cidr values (Heikki Linnakangas)

A patch in the December 2011 releases of PostgreSQL caused memory leakage in these operations, which could be significant in scenarios such as building a btree index on such a column.

• Avoid double close of file handle in syslogger on Windows (MauMau)

Ordinarily this error was invisible, but it would cause an exception when running on a debug version of Windows.

• Fix I/O-conversion-related memory leaks in plpgsql (Andres Freund, Jan Urbanski, Tom Lane)

Certain operations would leak memory until the end of the current function.

• Improve pg_dump's handling of inherited table columns (Tom Lane)

pg_dump mishandled situations where a child column has a different default expression than its parent column. If the default is textually identical to the parent's default, but not actually the same (for instance, because of schema search path differences) it would not be recognized as different, so that after dump and restore the child would be allowed to inherit the parent's default. Child columns that are NOT NULL where their parent is not could also be restored subtly incorrectly.

• Fix pg_restore's direct-to-database mode for INSERT-style table data (Tom Lane)

Direct-to-database restores from archive files made with --inserts or --column-inserts options fail when using pg_restore from a release dated September or December 2011, as a result of an oversight in a fix for another problem. The archive file itself is not at fault, and text-mode output is okay.

• Fix error in contrib/intarray's int[] & int[] operator (Guillaume Lelarge)
If the smallest integer the two input arrays have in common is 1, and there are smaller values in either array, then 1 would be incorrectly omitted from the result.

- Fix error detection in contrib/pgcrypto's encrypt_iv() and decrypt_iv() (Marko Kreen)

These functions failed to report certain types of invalid-input errors, and would instead return random garbage values for incorrect input.

- Fix one-byte buffer overrun in contrib/test_parser (Paul Guyot)

The code would try to read one more byte than it should, which would crash in corner cases. Since contrib/test_parser is only example code, this is not a security issue in itself, but bad example code is still bad.

- Use __sync_lock_test_and_set() for spinlocks on ARM, if available (Martin Pitt)

This function replaces our previous use of the SWPB instruction, which is deprecated and not available on ARMv6 and later. Reports suggest that the old code doesn't fail in an obvious way on recent ARM boards, but simply doesn't interlock concurrent accesses, leading to bizarre failures in multiprocess operation.

- Use -fexcess-precision=standard option when building with gcc versions that accept it (Andrew Dunstan)

This prevents assorted scenarios wherein recent versions of gcc will produce creative results.

- Allow use of threaded Python on FreeBSD (Chris Rees)

Our configure script previously believed that this combination wouldn't work; but FreeBSD fixed the problem, so remove that error check.

### E.188. Release 8.3.17

**Release date:** 2011-12-05

This release contains a variety of fixes from 8.3.16. For information about new features in the 8.3 major release, see Section E.205.

#### E.188.1. Migration to Version 8.3.17

A dump/restore is not required for those running 8.3.X.

However, a longstanding error was discovered in the definition of the information_schema.referential_constraints view. If you rely on correct results from that view, you should replace its definition as explained in the first changelog item below.

Also, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

#### E.188.2. Changes

- Fix bugs in information_schema.referential_constraints view (Tom Lane)

This view was being insufficiently careful about matching the foreign-key constraint to the depended-on primary or unique key constraint. That could result in failure to show a foreign key constraint at all, or showing it multiple times, or claiming that it depends on a different constraint than the one it really does.
Since the view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can (as a superuser) drop the `information_schema` schema then re-create it by sourcing `SHAREDIR/information_schema.sql`. (Run `pg_config --sharedir` if you're uncertain where `SHAREDIR` is.) This must be repeated in each database to be fixed.

- Fix TOAST-related data corruption during CREATE TABLE dest AS SELECT * FROM src or INSERT INTO dest SELECT * FROM src (Tom Lane)

  If a table has been modified by ALTER TABLE ADD COLUMN, attempts to copy its data verbatim to another table could produce corrupt results in certain corner cases. The problem can only manifest in this precise form in 8.4 and later, but we patched earlier versions as well in case there are other code paths that could trigger the same bug.

- Fix race condition during toast table access from stale syscache entries (Tom Lane)

  The typical symptom was transient errors like “missing chunk number 0 for toast value NNNNN in pg_toast_2619”, where the cited toast table would always belong to a system catalog.

- Make `DatumGetInetP()` unpack inet datums that have a 1-byte header, and add a new macro, `DatumGetInetPP()`, that does not (Heikki Linnakangas)

  This change affects no core code, but might prevent crashes in add-on code that expects `DatumGetInetP()` to produce an unpacked datum as per usual convention.

- Improve locale support in `money` type’s input and output (Tom Lane)

  Aside from not supporting all standard `lc_monetary` formatting options, the input and output functions were inconsistent, meaning there were locales in which dumped `money` values could not be re-read.

- Don’t let `transform_null_equals` affect CASE foo WHEN NULL ... constructs (Heikki Linnakangas)

  `transform_null_equals` is only supposed to affect `foo = NULL` expressions written directly by the user, not equality checks generated internally by this form of CASE.

- Change foreign-key trigger creation order to better support self-referential foreign keys (Tom Lane)

  For a cascading foreign key that references its own table, a row update will fire both the ON UPDATE trigger and the CHECK trigger as one event. The ON UPDATE trigger must execute first, else the CHECK will check a non-final state of the row and possibly throw an inappropriate error. However, the firing order of these triggers is determined by their names, which generally sort in creation order since the triggers have auto-generated names following the convention “RI_ConstraintTrigger_NNNN”. A proper fix would require modifying that convention, which we will do in 9.2, but it seems risky to change it in existing releases. So this patch just changes the creation order of the triggers. Users encountering this type of error should drop and re-create the foreign key constraint to get its triggers into the right order.

- Avoid floating-point underflow while tracking buffer allocation rate (Greg Matthews)

  While harmless in itself, on certain platforms this would result in annoying kernel log messages.

- Preserve blank lines within commands in psql’s command history (Robert Haas)

  The former behavior could cause problems if an empty line was removed from within a string literal, for example.

- Fix `pg_dump` to dump user-defined casts between auto-generated types, such as table rowtypes (Tom Lane)
• Use the preferred version of xsubpp to build PL/Perl, not necessarily the operating system's main copy (David Wheeler and Alex Hunsaker)

• Fix incorrect coding in contrib/dict_int and contrib/dict_xsyn (Tom Lane)
  Some functions incorrectly assumed that memory returned by palloc() is guaranteed zeroed.

• Honor query cancel interrupts promptly in pgstatindex() (Robert Haas)

• Ensure VPATH builds properly install all server header files (Peter Eisentraut)

• Shorten file names reported in verbose error messages (Peter Eisentraut)
  Regular builds have always reported just the name of the C file containing the error message call, but VPATH builds formerly reported an absolute path name.

• Fix interpretation of Windows timezone names for Central America (Tom Lane)
  Map “Central America Standard Time” to CST6, not CST6CDT, because DST is generally not observed anywhere in Central America.

• Update time zone data files to tzdata release 2011n for DST law changes in Brazil, Cuba, Fiji, Palestine, Russia, and Samoa; also historical corrections for Alaska and British East Africa.

E.189. Release 8.3.16

Release date: 2011-09-26

This release contains a variety of fixes from 8.3.15. For information about new features in the 8.3 major release, see Section E.205.

E.189.1. Migration to Version 8.3.16

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

E.189.2. Changes

• Fix bugs in indexing of in-doubt HOT-updated tuples (Tom Lane)
  These bugs could result in index corruption after reindexing a system catalog. They are not believed to affect user indexes.

• Fix multiple bugs in GiST index page split processing (Heikki Linnakangas)
  The probability of occurrence was low, but these could lead to index corruption.

• Fix possible buffer overrun in tsvector_concat() (Tom Lane)
  The function could underestimate the amount of memory needed for its result, leading to server crashes.

• Fix crash in xml_recv when processing a “standalone” parameter (Tom Lane)

• Avoid possibly accessing off the end of memory in ANALYZE and in SJIS-2004 encoding conversion (Noah Misch)
  This fixes some very-low-probability server crash scenarios.

• Fix race condition in relcache init file invalidation (Tom Lane)
There was a window wherein a new backend process could read a stale init file but miss the inval messages that would tell it the data is stale. The result would be bizarre failures in catalog accesses, typically “could not read block 0 in file ...” later during startup.

- Fix memory leak at end of a GiST index scan (Tom Lane)

Commands that perform many separate GiST index scans, such as verification of a new GiST-based exclusion constraint on a table already containing many rows, could transiently require large amounts of memory due to this leak.

- Fix performance problem when constructing a large, lossy bitmap (Tom Lane)

- Fix array- and path-creating functions to ensure padding bytes are zeroes (Tom Lane)

This avoids some situations where the planner will think that semantically-equal constants are not equal, resulting in poor optimization.

- Work around gcc 4.6.0 bug that breaks WAL replay (Tom Lane)

This could lead to loss of committed transactions after a server crash.

- Fix dump bug for VALUES in a view (Tom Lane)

- Disallow SELECT FOR UPDATE/SHARE on sequences (Tom Lane)

This operation doesn't work as expected and can lead to failures.

- Defend against integer overflow when computing size of a hash table (Tom Lane)

- Fix cases where CLUSTER might attempt to access already-removed TOAST data (Tom Lane)

- Fix portability bugs in use of credentials control messages for “peer” authentication (Tom Lane)

- Fix SSPI login when multiple roundtrips are required (Ahmed Shinwari, Magnus Hagander)

  The typical symptom of this problem was “The function requested is not supported” errors during SSPI login.

- Fix typo in pg_srand48 seed initialization (Andres Freund)

  This led to failure to use all bits of the provided seed. This function is not used on most platforms (only those without srand), and the potential security exposure from a less-random-than-expected seed seems minimal in any case.

- Avoid integer overflow when the sum of LIMIT and OFFSET values exceeds 2^63 (Heikki Linnakangas)

- Add overflow checks to int4 and int8 versions of generate_series() (Robert Haas)

- Fix trailing-zero removal in to_char() (Marti Raudsepp)

  In a format with FM and no digit positions after the decimal point, zeroes to the left of the decimal point could be removed incorrectly.

- Fix pg_size_pretty() to avoid overflow for inputs close to 2^63 (Tom Lane)

- In pg_ctl, support silent mode for service registrations on Windows (MauMau)

- Fix psql's counting of script file line numbers during COPY from a different file (Tom Lane)

- Fix pg_restore's direct-to-database mode for standard_conforming_strings (Tom Lane)
pg_restore could emit incorrect commands when restoring directly to a database server from an archive file that had been made with `standard_conforming_strings` set to on.

- Fix write-past-buffer-end and memory leak in libpq's LDAP service lookup code (Albe Laurenz)
- In libpq, avoid failures when using nonblocking I/O and an SSL connection (Martin Pihlak, Tom Lane)
- Improve libpq's handling of failures during connection startup (Tom Lane)

  In particular, the response to a server report of `fork()` failure during SSL connection startup is now saner.

- Improve libpq's error reporting for SSL failures (Tom Lane)
- Make ecpglib write double values with 15 digits precision (Akira Kurosawa)
- In ecpglib, be sure `LC_NUMERIC` setting is restored after an error (Michael Meskes)
- Apply upstream fix for blowfish signed-character bug (CVE-2011-2483) (Tom Lane)

  `contrib/pg_crypto`'s blowfish encryption code could give wrong results on platforms where `char` is signed (which is most), leading to encrypted passwords being weaker than they should be.

- Fix memory leak in `contrib/seg` (Heikki Linnakangas)
- Fix `pgstatindex()` to give consistent results for empty indexes (Tom Lane)
- Allow building with perl 5.14 (Alex Hunsaker)
- Update configure script's method for probing existence of system functions (Tom Lane)

  The version of autoconf we used in 8.3 and 8.2 could be fooled by compilers that perform link-time optimization.

- Fix assorted issues with build and install file paths containing spaces (Tom Lane)
- Update time zone data files to `tzdata` release 2011i for DST law changes in Canada, Egypt, Russia, Samoa, and South Sudan.

### E.190. Release 8.3.15

**Release date:** 2011-04-18

This release contains a variety of fixes from 8.3.14. For information about new features in the 8.3 major release, see Section E.205.

### E.190.1. Migration to Version 8.3.15

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

### E.190.2. Changes

- Disallow including a composite type in itself (Tom Lane)

  This prevents scenarios wherein the server could recurse infinitely while processing the composite type. While there are some possible uses for such a structure, they don't seem compelling enough to justify the effort required to make sure it always works safely.
• Avoid potential deadlock during catalog cache initialization (Nikhil Sontakke)

In some cases the cache loading code would acquire share lock on a system index before locking
the index's catalog. This could deadlock against processes trying to acquire exclusive locks in the
other, more standard order.

• Fix dangling-pointer problem in BEFORE ROW UPDATE trigger handling when there was a con-
current update to the target tuple (Tom Lane)

This bug has been observed to result in intermittent “cannot extract system attribute from virtual
tuple” failures while trying to do UPDATE RETURNING ctid. There is a very small probability
of more serious errors, such as generating incorrect index entries for the updated tuple.

• Disallow DROP TABLE when there are pending deferred trigger events for the table (Tom Lane)

Formerly the DROP would go through, leading to “could not open relation with OID nnn” errors
when the triggers were eventually fired.

• Fix PL/Python memory leak involving array slices (Daniel Popowich)

• Fix pg_restore to cope with long lines (over 1KB) in TOC files (Tom Lane)

• Put in more safeguards against crashing due to division-by-zero with overly enthusiastic compiler
optimization (Aurelien Jarno)

• Support use of dlopen() in FreeBSD and OpenBSD on MIPS (Tom Lane)

There was a hard-wired assumption that this system function was not available on MIPS hardware
on these systems. Use a compile-time test instead, since more recent versions have it.

• Fix compilation failures on HP-UX (Heikki Linnakangas)

• Fix version-incompatibility problem with libintl on Windows (Hiroshi Inoue)

• Fix usage of xcopy in Windows build scripts to work correctly under Windows 7 (Andrew Dunstan)

This affects the build scripts only, not installation or usage.

• Fix path separator used by pg_regress on Cygwin (Andrew Dunstan)

• Update time zone data files to tzdata release 2011f for DST law changes in Chile, Cuba, Falkland
Islands, Morocco, Samoa, and Turkey; also historical corrections for South Australia, Alaska, and
Hawaii.

E.191. Release 8.3.14

Release date: 2011-01-31

This release contains a variety of fixes from 8.3.13. For information about new features in the 8.3
major release, see Section E.205.

E.191.1. Migration to Version 8.3.14

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version
earlier than 8.3.8, see Section E.197.

E.191.2. Changes

• Avoid failures when EXPLAIN tries to display a simple-form CASE expression (Tom Lane)
If the CASE’s test expression was a constant, the planner could simplify the CASE into a form that confused the expression-display code, resulting in “unexpected CASE WHEN clause” errors.

- Fix assignment to an array slice that is before the existing range of subscripts (Tom Lane)
  
  If there was a gap between the newly added subscripts and the first pre-existing subscript, the code miscalculated how many entries needed to be copied from the old array’s null bitmap, potentially leading to data corruption or crash.

- Avoid unexpected conversion overflow in planner for very distant date values (Tom Lane)
  
  The date type supports a wider range of dates than can be represented by the timestamp types, but the planner assumed it could always convert a date to timestamp with impunity.

- Fix pg_restore’s text output for large objects (BLOBs) when standard_conforming_strings is on (Tom Lane)
  
  Although restoring directly to a database worked correctly, string escaping was incorrect if pg_restore was asked for SQL text output and standard_conforming_strings had been enabled in the source database.

- Fix erroneous parsing of tsquery values containing ... & !(subexpression) | ...

  (Tom Lane)
  
  Queries containing this combination of operators were not executed correctly. The same error existed in contrib/intarray’s query_int type and contrib/ltree’s ltxtquery type.

- Fix buffer overrun in contrib/intarray’s input function for the query_int type (Apple)
  
  This bug is a security risk since the function’s return address could be overwritten. Thanks to Apple Inc’s security team for reporting this issue and supplying the fix. (CVE-2010-4015)

- Fix bug in contrib/seg’s GiST picksplit algorithm (Alexander Korotkov)
  
  This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a seg column. If you have such an index, consider REINDEXing it after installing this update. (This is identical to the bug that was fixed in contrib/cube in the previous update.)

E.192. Release 8.3.13

Release date: 2010-12-16

This release contains a variety of fixes from 8.3.12. For information about new features in the 8.3 major release, see Section E.205.

E.192.1. Migration to Version 8.3.13

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

E.192.2. Changes

- Force the default wal_sync_method to be fdatasync on Linux (Tom Lane, Marti Raudsepp)
  
  The default on Linux has actually been fdatasync for many years, but recent kernel changes caused PostgreSQL to choose open_datasync instead. This choice did not result in any performance improvement, and caused outright failures on certain filesystems, notably ext4 with the data=journal mount option.
• Fix assorted bugs in WAL replay logic for GIN indexes (Tom Lane)

This could result in “bad buffer id: 0” failures or corruption of index contents during replication.

• Fix recovery from base backup when the starting checkpoint WAL record is not in the same WAL segment as its redo point (Jeff Davis)

• Fix persistent slowdown of autovacuum workers when multiple workers remain active for a long time (Tom Lane)

The effective vacuum_cost_limit for an autovacuum worker could drop to nearly zero if it processed enough tables, causing it to run extremely slowly.

• Add support for detecting register-stack overrun on IA64 (Tom Lane)

The IA64 architecture has two hardware stacks. Full prevention of stack-overrun failures requires checking both.

• Add a check for stack overflow in copyObject() (Tom Lane)

Certain code paths could crash due to stack overflow given a sufficiently complex query.

• Fix detection of page splits in temporary GiST indexes (Heikki Linnakangas)

It is possible to have a “concurrent” page split in a temporary index, if for example there is an open cursor scanning the index when an insertion is done. GiST failed to detect this case and hence could deliver wrong results when execution of the cursor continued.

• Avoid memory leakage while ANALYZE'ing complex index expressions (Tom Lane)

• Ensure an index that uses a whole-row Var still depends on its table (Tom Lane)

An index declared like create index i on t (foo(t.*) would not automatically get dropped when its table was dropped.

• Do not “inline” a SQL function with multiple OUT parameters (Tom Lane)

This avoids a possible crash due to loss of information about the expected result rowtype.

• Behave correctly if ORDER BY, LIMIT, FOR UPDATE, or WITH is attached to the VALUES part of INSERT ... VALUES (Tom Lane)

• Fix constant-folding of COALESCE() expressions (Tom Lane)

The planner would sometimes attempt to evaluate sub-expressions that in fact could never be reached, possibly leading to unexpected errors.

• Fix postmaster crash when connection acceptance (accept() or one of the calls made immediately after it) fails, and the postmaster was compiled with GSSAPI support (Alexander Chernikov)

• Fix missed unlink of temporary files when log_temp_files is active (Tom Lane)

If an error occurred while attempting to emit the log message, the unlink was not done, resulting in accumulation of temp files.

• Add print functionality for InhRelation nodes (Tom Lane)

This avoids a failure when debug_print_parse is enabled and certain types of query are executed.

• Fix incorrect calculation of distance from a point to a horizontal line segment (Tom Lane)
This bug affected several different geometric distance-measurement operators.

- Fix PL/pgSQL's handling of “simple” expressions to not fail in recursion or error-recovery cases (Tom Lane)

- Fix PL/Python's handling of set-returning functions (Jan Urbanski)

Attempts to call SPI functions within the iterator generating a set result would fail.

- Fix bug in contrib/cube's GiST picksplit algorithm (Alexander Korotkov)

This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a cube column. If you have such an index, consider REINDEXing it after installing this update.

- Don't emit “identifier will be truncated” notices in contrib/dblink except when creating new connections (Itagaki Takahiro)

- Fix potential coredump on missing public key in contrib/pgcrypto (Marti Raudsepp)

- Fix memory leak in contrib/xml2's XPath query functions (Tom Lane)

- Update time zone data files to tzdata release 2010o for DST law changes in Fiji and Samoa; also historical corrections for Hong Kong.

**E.193. Release 8.3.12**

**Release date:** 2010-10-04

This release contains a variety of fixes from 8.3.11. For information about new features in the 8.3 major release, see Section E.205.

**E.193.1. Migration to Version 8.3.12**

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

**E.193.2. Changes**

- Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function's owner.

The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).
• Prevent possible crashes in `pg_get_expr()` by disallowing it from being called with an argument that is not one of the system catalog columns it's intended to be used with (Heikki Linnakangas, Tom Lane)

• Treat exit code 128 (ERROR_WAIT_NO_CHILDREN) as non-fatal on Windows (Magnus Hagander)

Under high load, Windows processes will sometimes fail at startup with this error code. Formerly the postmaster treated this as a panic condition and restarted the whole database, but that seems to be an overreaction.

• Fix incorrect usage of non-strict OR join clauses in Append index scans (Tom Lane)

This is a back-patch of an 8.4 fix that was missed in the 8.3 branch. This corrects an error introduced in 8.3.8 that could cause incorrect results for outer joins when the inner relation is an inheritance tree or `UNION ALL` subquery.

• Fix possible duplicate scans of `UNION ALL` member relations (Tom Lane)

• Fix “cannot handle unplanned sub-select” error (Tom Lane)

This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.

• Fix failure to mark cached plans as transient (Tom Lane)

If a plan is prepared while `CREATE INDEX CONCURRENTLY` is in progress for one of the referenced tables, it is supposed to be re-planned once the index is ready for use. This was not happening reliably.

• Reduce PANIC to ERROR in some occasionally-reported btree failure cases, and provide additional detail in the resulting error messages (Tom Lane)

This should improve the system's robustness with corrupted indexes.

• Prevent show_session_authorization() from crashing within autovacuum processes (Tom Lane)

• Defend against functions returning setof record where not all the returned rows are actually of the same rowtype (Tom Lane)

• Fix possible failure when hashing a pass-by-reference function result (Tao Ma, Tom Lane)

• Improve merge join's handling of NULLs in the join columns (Tom Lane)

A merge join can now stop entirely upon reaching the first NULL, if the sort order is such that NULLs sort high.

• Take care to fsync the contents of lockfiles (both `postmaster.pid` and the socket lockfile) while writing them (Tom Lane)

This omission could result in corrupted lockfile contents if the machine crashes shortly after postmaster start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.

• Avoid recursion while assigning XIDs to heavily-nested subtransactions (Andres Freund, Robert Haas)

The original coding could result in a crash if there was limited stack space.

• Avoid holding open old WAL segments in the walwriter process (Magnus Hagander, Heikki Linnakangas)
The previous coding would prevent removal of no-longer-needed segments.

- Fix log_line_prefix's %i escape, which could produce junk early in backend startup (Tom Lane)
- Fix possible data corruption in ALTER TABLE ... SET TABLESPACE when archiving is enabled (Jeff Davis)
- Allow CREATE DATABASE and ALTER DATABASE ... SET TABLESPACE to be interrupted by query-cancel (Guillaume Lelarge)
- Fix REASSIGN OWNED to handle operator classes and families (Asko Tiidumaa)
- Fix possible core dump when comparing two empty tsquery values (Tom Lane)
- Fix LIKE's handling of patterns containing % followed by _ (Tom Lane)

We've fixed this before, but there were still some incorrectly-handled cases.

- In PL/Python, defend against null pointer results from PyCObject_AsVoidPtr and PyCObject_FromVoidPtr (Peter Eisentraut)
- Make psql recognize DISCARD ALL as a command that should not be encased in a transaction block in autocommit-off mode (Itagaki Takahiro)
- Fix ecpg to process data from RETURNING clauses correctly (Michael Meskes)
- Improve contrib/dblink's handling of tables containing dropped columns (Tom Lane)
- Fix connection leak after “duplicate connection name” errors in contrib/dblink (Itagaki Takahiro)
- Fix contrib/dblink to handle connection names longer than 62 bytes correctly (Itagaki Takahiro)
- Add hstore(text, text) function to contrib/hstore (Robert Haas)

This function is the recommended substitute for the now-deprecated => operator. It was back-patched so that future-proofed code can be used with older server versions. Note that the patch will be effective only after contrib/hstore is installed or reinstalled in a particular database. Users might prefer to execute the CREATE FUNCTION command by hand, instead.

- Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)
- Update time zone data files to tzdata release 2010l for DST law changes in Egypt and Palestine; also historical corrections for Finland.

This change also adds new names for two Micronesian timezones: Pacific/Chuuk is now preferred over Pacific/Truk (and the preferred abbreviation is CHUT not TRUT) and Pacific/Pohnpei is preferred over Pacific/Ponape.

- Make Windows' “N. Central Asia Standard Time” timezone map to Asia/Novosibirsk, not Asia/Almaty (Magnus Hagander)

Microsoft changed the DST behavior of this zone in the timezone update from KB976098. Asia/Novosibirsk is a better match to its new behavior.

E.194. Release 8.3.11

Release date: 2010-05-17
This release contains a variety of fixes from 8.3.10. For information about new features in the 8.3 major release, see Section E.205.

E.194.1. Migration to Version 8.3.11

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

E.194.2. Changes

- Enforce restrictions in plperl using an opmask applied to the whole interpreter, instead of using Safe.pm (Tim Bunce, Andrew Dunstan)

  Recent developments have convinced us that Safe.pm is too insecure to rely on for making plperl trustworthy. This change removes use of Safe.pm altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include that it is now possible to use Perl's strict pragma in a natural way in plperl, and that Perl's @a and $b variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

- Prevent PL/Tcl from executing untrustworthy code from pltcl_modules (Tom)

  PL/Tcl's feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless pltcl_modules is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted “normal” Tcl interpreter unless we are really going to execute a pltclu function. (CVE-2010-1170)

- Fix possible crash if a cache reset message is received during rebuild of a relcache entry (Heikki)

  This error was introduced in 8.3.10 while fixing a related failure.

- Apply per-function GUC settings while running the language validator for the function (Itagaki Takahiro)

  This avoids failures if the function's code is invalid without the setting; an example is that SQL functions may not parse if the search_path is not correct.

- Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

  Previously, if an unprivileged user ran ALTER USER ... RESET ALL for himself, or ALTER DATABASE ... RESET ALL for a database he owns, this would remove all special parameter settings for the user or database, even ones that are only supposed to be changeable by a superuser. Now, the ALTER will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a CONTEXT addition would be made to log entries (Tom)

  In some cases the context-printing function would fail because the current transaction had already been rolled back when it came time to print a log message.

- Ensure the archiver process responds to changes in archive_command as soon as possible (Tom)

- Update PL/Perl's ppport.h for modern Perl versions (Andrew)

- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)

- Prevent infinite recursion in psql when expanding a variable that refers to itself (Tom)

3003
• Fix psql's `\copy` to not add spaces around a dot within `\copy (select ...)` (Tom)

   Addition of spaces around the decimal point in a numeric literal would result in a syntax error.

• Fix unnecessary “GIN indexes do not support whole-index scans” errors for unsatisfiable queries using `contrib/intarray` operators (Tom)

• Ensure that `contrib/pgstattuple` functions respond to cancel interrupts promptly (Tatsuhito Kasahara)

• Make server startup deal properly with the case that `shmget()` returns `EINVAL` for an existing shared memory segment (Tom)

   This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entire-
   ly-misleading startup failure complaining that the shared memory request size was too large.

• Avoid possible crashes in syslogger process on Windows (Heikki)

• Deal more robustly with incomplete time zone information in the Windows registry (Magnus)

• Update the set of known Windows time zone names (Magnus)

• Update time zone data files to tzdata release 2010j for DST law changes in Argentina, Australian
   Antarctic, Bangladesh, Mexico, Morocco, Pakistan, Palestine, Russia, Syria, Tunisia; also historical
   corrections for Taiwan.

   Also, add `PKST` (Pakistan Summer Time) to the default set of timezone abbreviations.

E.195. Release 8.3.10

Release date: 2010-03-15

This release contains a variety of fixes from 8.3.9. For information about new features in the 8.3 major
release, see Section E.205.

E.195.1. Migration to Version 8.3.10

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version
earlier than 8.3.8, see Section E.197.

E.195.2. Changes

• Add new configuration parameter `ssl_renegotiation_limit` to control how often we do
  session key renegotiation for an SSL connection (Magnus)

   This can be set to zero to disable renegotiation completely, which may be required if a broken SSL
   library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that
   cause renegotiation attempts to fail.

• Fix possible deadlock during backend startup (Tom)

• Fix possible crashes due to not handling errors during relcache reload cleanly (Tom)

• Fix possible crash due to use of dangling pointer to a cached plan (Tatsuo)

• Fix possible crashes when trying to recover from a failure in subtransaction start (Tom)

• Fix server memory leak associated with use of savepoints and a client encoding different from
  server's encoding (Tom)
• Fix incorrect WAL data emitted during end-of-recovery cleanup of a GIST index page split (Yoichi Hirai)
  
  This would result in index corruption, or even more likely an error during WAL replay, if we were unlucky enough to crash during end-of-recovery cleanup after having completed an incomplete GIST insertion.

• Make substring() for bit types treat any negative length as meaning “all the rest of the string” (Tom)
  
  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

• Fix integer-to-bit-string conversions to handle the first fractional byte correctly when the output bit width is wider than the given integer by something other than a multiple of 8 bits (Tom)

• Fix some cases of pathologically slow regular expression matching (Tom)

• Fix assorted crashes in xml processing caused by sloppy memory management (Tom)
  
  This is a back-patch of changes first applied in 8.4. The 8.3 code was known buggy, but the new code was sufficiently different to not want to back-patch it until it had gotten some field testing.

• Fix bug with trying to update a field of an element of a composite-type array column (Tom)

• Fix the STOP WAL LOCATION entry in backup history files to report the next WAL segment's name when the end location is exactly at a segment boundary (Itagaki Takahiro)

• Fix some more cases of temporary-file leakage (Heikki)
  
  This corrects a problem introduced in the previous minor release. One case that failed is when a plpgsql function returning set is called within another function's exception handler.

• Improve constraint exclusion processing of boolean-variable cases, in particular make it possible to exclude a partition that has a “bool_column = false” constraint (Tom)

• When reading pg_hba.conf and related files, do not treat @something as a file inclusion request if the @ appears inside quote marks; also, never treat @ by itself as a file inclusion request (Tom)
  
  This prevents erratic behavior if a role or database name starts with @. If you need to include a file whose path name contains spaces, you can still do so, but you must write @"/path to/file" rather than putting the quotes around the whole construct.

• Prevent infinite loop on some platforms if a directory is named as an inclusion target in pg_hba.conf and related files (Tom)

• Fix possible infinite loop if SSL_read or SSL_write fails without setting errno (Tom)
  
  This is reportedly possible with some Windows versions of OpenSSL.

• Disallow GSSAPI authentication on local connections, since it requires a hostname to function correctly (Magnus)

• Make ecpg report the proper SQLSTATE if the connection disappears (Michael)

• Fix psql's numericlocale option to not format strings it shouldn’t in latex and troff output formats (Heikki)

• Make psql return the correct exit status (3) when ON_ERROR_STOP and --single-transaction are both specified and an error occurs during the implied COMMIT (Bruce)
• Fix plpgsql failure in one case where a composite column is set to NULL (Tom)
• Fix possible failure when calling PL/Perl functions from PL/PerlU or vice versa (Tim Bunce)
• Add volatile markings in PL/Python to avoid possible compiler-specific misbehavior (Zdenek Kotala)
• Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)
  The only known symptom of this oversight is that the Tcl clock command misbehaves if using Tcl 8.5 or later.
• Prevent crash in contrib/dblink when too many key columns are specified to a dblink_build_sql_* function (Rushabh Lathia, Joe Conway)
• Allow zero-dimensional arrays in contrib/ltree operations (Tom)
  This case was formerly rejected as an error, but it's more convenient to treat it the same as a zero-element array. In particular this avoids unnecessary failures when an ltree operation is applied to the result of ARRAY(SELECT ...) and the sub-select returns no rows.
• Fix assorted crashes in contrib/xml2 caused by sloppy memory management (Tom)
• Make building of contrib/xml2 more robust on Windows (Andrew)
• Fix race condition in Windows signal handling (Radu Ilie)
  One known symptom of this bug is that rows in pg_listener could be dropped under heavy load.
• Update time zone data files to tzdata release 2010e for DST law changes in Bangladesh, Chile, Fiji, Mexico, Paraguay, Samoa.

E.196. Release 8.3.9

Release date: 2009-12-14

This release contains a variety of fixes from 8.3.8. For information about new features in the 8.3 major release, see Section E.205.

E.196.1. Migration to Version 8.3.9

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.8, see Section E.197.

E.196.2. Changes

• Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)
  This change prevents allegedly-immutable index functions from possibly subverting a superuser's session (CVE-2009-4136).
• Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)
  This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).
• Fix possible crash during backend-startup-time cache initialization (Tom)
• Avoid crash on empty thesaurus dictionary (Tom)
- Prevent signals from interrupting VACUUM at unsafe times (Alvaro)

  This fix prevents a PANIC if a VACUUM FULL is canceled after it's already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

- Fix possible crash due to integer overflow in hash table size calculation (Tom)

  This could occur with extremely large planner estimates for the size of a hashjoin's result.

- Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)

- Ensure that shared tuple-level locks held by prepared transactions are not ignored (Heikki)

- Fix premature drop of temporary files used for a cursor that is accessed within a subtransaction (Heikki)

- Fix memory leak in syslogger process when rotating to a new CSV logfile (Tom)

- Fix Windows permission-downgrade logic (Jesse Morris)

  This fixes some cases where the database failed to start on Windows, often with misleading error messages such as “could not locate matching postgres executable”.

- Fix incorrect logic for GiST index page splits, when the split depends on a non-first column of the index (Paul Ramsey)

- Don't error out if recycling or removing an old WAL file fails at the end of checkpoint (Heikki)

  It's better to treat the problem as non-fatal and allow the checkpoint to complete. Future checkpoints will retry the removal. Such problems are not expected in normal operation, but have been seen to be caused by misdesigned Windows anti-virus and backup software.

- Ensure WAL files aren't repeatedly archived on Windows (Heikki)

  This is another symptom that could happen if some other process interfered with deletion of a no-longer-needed file.

- Fix PAM password processing to be more robust (Tom)

  The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.

- Raise the maximum authentication token (Kerberos ticket) size in GSSAPI and SSPI authentication methods (Ian Turner)

  While the old 2000-byte limit was more than enough for Unix Kerberos implementations, tickets issued by Windows Domain Controllers can be much larger.

- Re-enable collection of access statistics for sequences (Akira Kurosawa)

  This used to work but was broken in 8.3.

- Fix processing of ownership dependencies during CREATE OR REPLACE FUNCTION (Tom)

- Fix incorrect handling of WHERE \( x=x \) conditions (Tom)

  In some cases these could get ignored as redundant, but they aren't — they're equivalent to \( x \text{ IS NOT NULL} \).

- Make text search parser accept underscores in XML attributes (Peter)
• Fix encoding handling in xml binary input (Heikki)
  If the XML header doesn't specify an encoding, we now assume UTF-8 by default; the previous handling was inconsistent.

• Fix bug with calling plperl from plperlu or vice versa (Tom)
  An error exit from the inner function could result in crashes due to failure to re-select the correct Perl interpreter for the outer function.

• Fix session-lifespan memory leak when a PL/Perl function is redefined (Tom)

• Ensure that Perl arrays are properly converted to PostgreSQL arrays when returned by a set-returning PL/Perl function (Andrew Dunstan, Abhijit Menon-Sen)
  This worked correctly already for non-set-returning functions.

• Fix rare crash in exception processing in PL/Python (Peter)

• In contrib/pg_standby, disable triggering failover with a signal on Windows (Fujii Masao)
  This never did anything useful, because Windows doesn't have Unix-style signals, but recent changes made it actually crash.

• Ensure psql's flex module is compiled with the correct system header definitions (Tom)
  This fixes build failures on platforms where --enable-largefile causes incompatible changes in the generated code.

• Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

• Update the timezone abbreviation files to match current reality (Joachim Wieland)
  This includes adding IDT and SGT to the default timezone abbreviation set.

• Update time zone data files to tzdata release 2009s for DST law changes in Antarctica, Argentina, Bangladesh, Fiji, Novokuznetsk, Pakistan, Palestine, Samoa, Syria; also historical corrections for Hong Kong.

E.197. Release 8.3.8

Release date: 2009-09-09

This release contains a variety of fixes from 8.3.7. For information about new features in the 8.3 major release, see Section E.205.

E.197.1. Migration to Version 8.3.8

A dump/restore is not required for those running 8.3.X. However, if you have any hash indexes on interval columns, you must REINDEX them after updating to 8.3.8. Also, if you are upgrading from a version earlier than 8.3.5, see Section E.200.

E.197.2. Changes

• Fix Windows shared-memory allocation code (Tsutomu Yamada, Magnus)
  This bug led to the often-reported “could not reattach to shared memory” error message.

• Force WAL segment switch during pg_start_backup() (Heikki)
This avoids corner cases that could render a base backup unusable.

- Disallow \texttt{RESET ROLE} and \texttt{RESET SESSION AUTHORIZATION} inside security-definer functions (Tom, Heikki)

  This covers a case that was missed in the previous patch that disallowed \texttt{SET ROLE} and \texttt{SET SESSION AUTHORIZATION} inside security-definer functions. (See CVE-2007-6600)

- Make \texttt{LOAD} of an already-loaded loadable module into a no-op (Tom)

  Formerly, \texttt{LOAD} would attempt to unload and re-load the module, but this is unsafe and not all that useful.

- Disallow empty passwords during LDAP authentication (Magnus)

- Fix handling of sub-\texttt{SELECT}s appearing in the arguments of an outer-level aggregate function (Tom)

- Fix bugs associated with fetching a whole-row value from the output of a Sort or Materialize plan node (Tom)

- Prevent \texttt{synchronize_seqscans} from changing the results of scrollable and \texttt{WITH HOLD} cursors (Tom)

- Revert planner change that disabled partial-index and constraint exclusion optimizations when there were more than 100 clauses in an \texttt{AND} or \texttt{OR} list (Tom)

- Fix hash calculation for data type \texttt{interval} (Tom)

  This corrects wrong results for hash joins on interval values. It also changes the contents of hash indexes on interval columns. If you have any such indexes, you must \texttt{REINDEX} them after updating.

- Treat \texttt{to_char(..., 'TH')} as an uppercase ordinal suffix with 'HH'/'HH12' (Heikki)

  It was previously handled as 'th' (lowercase).

- Fix overflow for \texttt{INTERVAL 'x ms'} when \texttt{x} is more than 2 million and integer datetimes are in use (Alex Hunsaker)

- Fix calculation of distance between a point and a line segment (Tom)

  This led to incorrect results from a number of geometric operators.

- Fix \texttt{money} data type to work in locales where currency amounts have no fractional digits, e.g. Japan (Itagaki Takahiro)

- Fix \texttt{LIKE} for case where pattern contains \texttt{%_} (Tom)

- Properly round datetime input like \texttt{00:12:57.999999999999999999999999} (Tom)

- Fix memory leaks in XML operations (Tom)

- Fix poor choice of page split point in GiST R-tree operator classes (Teodor)

- Ensure that a “fast shutdown” request will forcibly terminate open sessions, even if a “smart shutdown” was already in progress (Fujii Masao)

- Avoid performance degradation in bulk inserts into GIN indexes when the input values are (nearly) in sorted order (Tom)

- Correctly enforce NOT NULL domain constraints in some contexts in PL/pgSQL (Tom)
• Fix portability issues in plperl initialization (Andrew Dunstan)
• Fix pg_ctl to not go into an infinite loop ifpostgresql.conf is empty (Jeff Davis)
• Improve pg_dump's efficiency when there are many large objects (Tamas Vincze)
• Use SIGUSR1, not SIGQUIT, as the failover signal for pg_standby (Heikki)
• Make pg_standby's maxretries option behave as documented (Fujii Masao)
• Make contrib/hstore throw an error when a key or value is too long to fit in its data structure, rather than silently truncating it (Andrew Gierth)
• Fix contrib/xml2’s xslt_process() to properly handle the maximum number of parameters (twenty) (Tom)
• Improve robustness of libpq's code to recover from errors during COPY FROM STDIN (Tom)
• Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)
• Update time zone data files to tzdata release 2009l for DST law changes in Bangladesh, Egypt, Jordan, Pakistan, Argentina/San_Luis, Cuba, Jordan (historical correction only), Mauritius, Morocco, Palestine, Syria, Tunisia.

E.198. Release 8.3.7

Release date: 2009-03-16

This release contains a variety of fixes from 8.3.6. For information about new features in the 8.3 major release, see Section E.205.

E.198.1. Migration to Version 8.3.7

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.5, see Section E.200.

E.198.2. Changes

• Prevent error recursion crashes when encoding conversion fails (Tom)

This change extends fixes made in the last two minor releases for related failure scenarios. The previous fixes were narrowly tailored for the original problem reports, but we have now recognized that any error thrown by an encoding conversion function could potentially lead to infinite recursion while trying to report the error. The solution therefore is to disable translation and encoding conversion and report the plain-ASCII form of any error message, if we find we have gotten into a recursive error reporting situation. (CVE-2009-0922)

• Disallow CREATE CONVERSION with the wrong encodings for the specified conversion function (Heikki)

This prevents one possible scenario for encoding conversion failure. The previous change is a backstop to guard against other kinds of failures in the same area.

• Fix xpath() to not modify the path expression unless necessary, and to make a saner attempt at it when necessary (Andrew)

The SQL standard suggests that xpath should work on data that is a document fragment, but libxml doesn't support that, and indeed it's not clear that this is sensible according to the XPath
xpath attempted to work around this mismatch by modifying both the data and the path expression, but the modification was buggy and could cause valid searches to fail. Now, xpath checks whether the data is in fact a well-formed document, and if so invokes libxml with no change to the data or path expression. Otherwise, a different modification method that is somewhat less likely to fail is used.

**Note**

The new modification method is still not 100% satisfactory, and it seems likely that no real solution is possible. This patch should therefore be viewed as a band-aid to keep from breaking existing applications unnecessarily. It is likely that PostgreSQL 8.4 will simply reject use of xpath on data that is not a well-formed document.

- Fix core dump when to_char() is given format codes that are inappropriate for the type of the data argument (Tom)
- Fix possible failure in text search when C locale is used with a multi-byte encoding (Teodor)
  
  Crashes were possible on platforms where wchar_t is narrower than int; Windows in particular.
- Fix extreme inefficiency in text search parser's handling of an email-like string containing multiple @ characters (Heikki)
- Fix planner problem with sub-SELECT in the output list of a larger subquery (Tom)
  
  The known symptom of this bug is a “failed to locate grouping columns” error that is dependent on the datatype involved; but there could be other issues as well.
- Fix decompilation of CASE WHEN with an implicit coercion (Tom)
  
  This mistake could lead to Assert failures in an Assert-enabled build, or an “unexpected CASE WHEN clause” error message in other cases, when trying to examine or dump a view.
- Fix possible misassignment of the owner of a TOAST table's rowtype (Tom)
  
  If CLUSTER or a rewriting variant of ALTER TABLE were executed by someone other than the table owner, the pg_type entry for the table's TOAST table would end up marked as owned by that someone. This caused no immediate problems, since the permissions on the TOAST rowtype aren't examined by any ordinary database operation. However, it could lead to unexpected failures if one later tried to drop the role that issued the command (in 8.1 or 8.2), or “owner of data type appears to be invalid” warnings from pg_dump after having done so (in 8.3).
- Change UNLISTEN to exit quickly if the current session has never executed any LISTEN command (Tom)
  
  Most of the time this is not a particularly useful optimization, but since DISCARD ALL invokes UNLISTEN, the previous coding caused a substantial performance problem for applications that made heavy use of DISCARD ALL.
- Fix PL/pgSQL to not treat INTO after INSERT as an INTO-variables clause anywhere in the string, not only at the start; in particular, don't fail for INSERT INTO within CREATE RULE (Tom)
- Clean up PL/pgSQL error status variables fully at block exit (Ashesh Vashi and Dave Page)
  
  This is not a problem for PL/pgSQL itself, but the omission could cause the PL/pgSQL Debugger to crash while examining the state of a function.
- Retry failed calls to CallNamedPipe() on Windows (Steve Marshall, Magnus)
It appears that this function can sometimes fail transiently; we previously treated any failure as a hard error, which could confuse LISTEN/NOTIFY as well as other operations.

- Add **MUST** (Mauritius Island Summer Time) to the default list of known timezone abbreviations (Xavier Bugaud)

**E.199. Release 8.3.6**

**Release date:** 2009-02-02

This release contains a variety of fixes from 8.3.5. For information about new features in the 8.3 major release, see Section E.205.

**E.199.1. Migration to Version 8.3.6**

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.5, see Section E.200.

**E.199.2. Changes**

- **Make DISCARD ALL release advisory locks, in addition to everything it already did** (Tom)

  This was decided to be the most appropriate behavior. This could affect existing applications, however.

- **Fix whole-index GiST scans to work correctly** (Teodor)

  This error could cause rows to be lost if a table is clustered on a GiST index.

- **Fix crash of xmlconcat(NULL)** (Peter)

- **Fix possible crash in ispell dictionary if high-bit-set characters are used as flags** (Teodor)

  This is known to be done by one widely available Norwegian dictionary, and the same condition may exist in others.

- **Fix misordering of pg_dump output for composite types** (Tom)

  The most likely problem was for user-defined operator classes to be dumped after indexes or views that needed them.

- **Improve handling of URLs in headline() function** (Teodor)

- **Improve handling of overlength headlines in headline() function** (Teodor)

- **Prevent possible Assert failure or misconversion if an encoding conversion is created with the wrong conversion function for the specified pair of encodings** (Tom, Heikki)

- **Fix possible Assert failure if a statement executed in PL/pgSQL is rewritten into another kind of statement, for example if an INSERT is rewritten into an UPDATE** (Heikki)

- **Ensure that a snapshot is available to datatype input functions** (Tom)

  This primarily affects domains that are declared with CHECK constraints involving user-defined stable or immutable functions. Such functions typically fail if no snapshot has been set.

- **Make it safer for SPI-using functions to be used within datatype I/O; in particular, to be used in domain check constraints** (Tom)
• Avoid unnecessary locking of small tables in VACUUM (Heikki)

• Fix a problem that sometimes kept ALTER TABLE ENABLE/DISABLE RULE from being recognized by active sessions (Tom)

• Fix a problem that made UPDATE RETURNING tableoid return zero instead of the correct OID (Tom)

• Allow functions declared as taking ANYARRAY to work on the pg_statistic columns of that type (Tom)

This used to work, but was unintentionally broken in 8.3.

• Fix planner misestimation of selectivity when transitive equality is applied to an outer-join clause (Tom)

This could result in bad plans for queries like ... from a left join b on a.a1 = b.b1 where a.a1 = 42 ...

• Improve optimizer's handling of long IN lists (Tom)

This change avoids wasting large amounts of time on such lists when constraint exclusion is enabled.

• Prevent synchronous scan during GIN index build (Tom)

Because GIN is optimized for inserting tuples in increasing TID order, choosing to use a synchronous scan could slow the build by a factor of three or more.

• Ensure that the contents of a holdable cursor don't depend on the contents of TOAST tables (Tom)

Previously, large field values in a cursor result might be represented as TOAST pointers, which would fail if the referenced table got dropped before the cursor is read, or if the large value is deleted and then vacuumed away. This cannot happen with an ordinary cursor, but it could with a cursor that is held past its creating transaction.

• Fix memory leak when a set-returning function is terminated without reading its whole result (Tom)

• Fix encoding conversion problems in XML functions when the database encoding isn't UTF-8 (Tom)

• Fix contrib/dblink's dblink_get_result(text,bool) function (Joe)

• Fix possible garbage output from contrib/sslinfo functions (Tom)

• Fix incorrect behavior of contrib/tsearch2 compatibility trigger when it's fired more than once in a command (Teodor)

• Fix possible mis-signaling in autovacuum (Heikki)

• Support running as a service on Windows 7 beta (Dave and Magnus)

• Fix ecpg's handling of varchar structs (Michael)

• Fix configure script to properly report failure when unable to obtain linkage information for PL/Perl (Andrew)

• Make all documentation reference pgsql-bugs and/or pgsql-hackers as appropriate, instead of the now-decommissioned pgsql-ports and pgsql-patches mailing lists (Tom)

• Update time zone data files to tzdata release 2009a (for Kathmandu and historical DST corrections in Switzerland, Cuba)
E.200. Release 8.3.5

Release date: 2008-11-03

This release contains a variety of fixes from 8.3.4. For information about new features in the 8.3 major release, see Section E.205.

E.200.1. Migration to Version 8.3.5

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.1, see Section E.204. Also, if you were running a previous 8.3.X release, it is recommended to REINDEX all GiST indexes after the upgrade.

E.200.2. Changes

• Fix GiST index corruption due to marking the wrong index entry “dead” after a deletion (Teodor)
  This would result in index searches failing to find rows they should have found. Corrupted indexes can be fixed with REINDEX.

• Fix backend crash when the client encoding cannot represent a localized error message (Tom)
  We have addressed similar issues before, but it would still fail if the “character has no equivalent” message itself couldn't be converted. The fix is to disable localization and send the plain ASCII error message when we detect such a situation.

• Fix possible crash in bytea-to-XML mapping (Michael McMaster)

• Fix possible crash when deeply nested functions are invoked from a trigger (Tom)

• Improve optimization of expression IN(expression-list) queries (Tom, per an idea from Robert Haas)
  Cases in which there are query variables on the right-hand side had been handled less efficiently in 8.2.x and 8.3.x than in prior versions. The fix restores 8.1 behavior for such cases.

• Fix mis-expansion of rule queries when a sub-SELECT appears in a function call in FROM, a multi-row VALUES list, or a RETURNING list (Tom)
  The usual symptom of this problem is an “unrecognized node type” error.

• Fix Assert failure during rescan of an IS NULL search of a GiST index (Teodor)

• Fix memory leak during rescan of a hashed aggregation plan (Neil)

• Ensure an error is reported when a newly-defined PL/pgSQL trigger function is invoked as a normal function (Tom)

• Force a checkpoint before CREATE DATABASE starts to copy files (Heikki)
  This prevents a possible failure if files had recently been deleted in the source database.

• Prevent possible collision of relfilenode numbers when moving a table to another tablespace with ALTER SET TABLESPACE (Heikki)
  The command tried to re-use the existing filename, instead of picking one that is known unused in the destination directory.

• Fix incorrect text search headline generation when single query item matches first word of text (Sushant Sinha)
• Fix improper display of fractional seconds in interval values when using a non-ISO datestyle in an
  --enable-integer-datetimes build (Ron Mayer)
• Make ILIKE compare characters case-insensitively even when they're escaped (Andrew)
• Ensure DISCARD is handled properly by statement logging (Tom)
• Fix incorrect logging of last-completed-transaction time during PITR recovery (Tom)
• Ensure SPI_getvalue and SPI_getbinval behave correctly when the passed tuple and tuple
descriptor have different numbers of columns (Tom)
  This situation is normal when a table has had columns added or removed, but these two functions
didn't handle it properly. The only likely consequence is an incorrect error indication.
• Mark SessionReplicationRole as PGDLLIMPORT so it can be used by Slony on Windows
  (Magnus)
• Fix small memory leak when using libpq's gsslib parameter (Magnus)
  The space used by the parameter string was not freed at connection close.
• Ensure libgssapi is linked into libpq if needed (Markus Schaaf)
• Fix ecpg's parsing of CREATE ROLE (Michael)
• Fix recent breakage of pg_ctl restart (Tom)
  pg_controldata and pg_resetxlog did this incorrectly, and so could fail on Windows.
• Update time zone data files to tzdata release 2008i (for DST law changes in Argentina, Brazil,
  Mauritius, Syria)

E.201. Release 8.3.4

Release date: 2008-09-22

This release contains a variety of fixes from 8.3.3. For information about new features in the 8.3 major
release, see Section E.205.

E.201.1. Migration to Version 8.3.4

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version
earlier than 8.3.1, see Section E.204.

E.201.2. Changes

• Fix bug in btree WAL recovery code (Heikki)
  Recovery failed if the WAL ended partway through a page split operation.
• Fix potential use of wrong cutoff XID for HOT page pruning (Alvaro)
  This error created a risk of corruption in system catalogs that are consulted by VACUUM: dead tuple
  versions might be removed too soon. The impact of this on actual database operations would be
  minimal, since the system doesn't follow MVCC rules while examining catalogs, but it might result
  in transiently wrong output from pg_dump or other client programs.
• Fix potential miscalculation of datfrozenxid (Alvaro)
This error may explain some recent reports of failure to remove old `pg_clog` data.

- Fix incorrect HOT updates after `pg_class` is reindexed (Tom)
  
  Corruption of `pg_class` could occur if `REINDEX TABLE pg_class` was followed in the same session by an `ALTER TABLE RENAME` or `ALTER TABLE SET SCHEMA` command.

- Fix missed “combo cid” case (Karl Schnaitter)

  This error made rows incorrectly invisible to a transaction in which they had been deleted by multiple subtransactions that all aborted.

- Prevent autovacuum from crashing if the table it's currently checking is deleted at just the wrong time (Alvaro)

- Widen local lock counters from 32 to 64 bits (Tom)

  This responds to reports that the counters could overflow in sufficiently long transactions, leading to unexpected “lock is already held” errors.

- Fix missed permissions checks when a view contains a simple `UNION ALL` construct (Heikki)

  Permissions for the referenced tables were checked properly, but not permissions for the view itself.

- Add checks in executor startup to ensure that the tuples produced by an `INSERT` or `UPDATE` will match the target table's current rowtype (Tom)

  This situation is believed to be impossible in 8.3, but it can happen in prior releases, so a check seems prudent.

- Fix possible repeated drops during `DROP OWNED` (Tom)

  This would typically result in strange errors such as “cache lookup failed for relation NNN”.

- Fix several memory leaks in XML operations (Kris Jurka, Tom)

- Fix `xmlserialize()` to raise error properly for unacceptable target data type (Tom)

- Fix a couple of places that mis-handled multibyte characters in text search configuration file parsing (Tom)

  Certain characters occurring in configuration files would always cause “invalid byte sequence for encoding” failures.

- Provide file name and line number location for all errors reported in text search configuration files (Tom)

- Fix `AT TIME ZONE` to first try to interpret its timezone argument as a timezone abbreviation, and only try it as a full timezone name if that fails, rather than the other way around as formerly (Tom)

  The timestamp input functions have always resolved ambiguous zone names in this order. Making `AT TIME ZONE` do so as well improves consistency, and fixes a compatibility bug introduced in 8.1: in ambiguous cases we now behave the same as 8.0 and before did, since in the older versions `AT TIME ZONE` accepted only abbreviations.
• Fix datetime input functions to correctly detect integer overflow when running on a 64-bit platform (Tom)

• Prevent integer overflows during units conversion when displaying a configuration parameter that has units (Tom)

• Improve performance of writing very long log messages to syslog (Tom)

• Allow spaces in the suffix part of an LDAP URL in pg_hba.conf (Tom)

• Fix bug in backwards scanning of a cursor on a SELECT DISTINCT ON query (Tom)

• Fix planner bug that could improperly push down IS NULL tests below an outer join (Tom)
  This was triggered by occurrence of IS NULL tests for the same relation in all arms of an upper OR clause.

• Fix planner bug with nested sub-select expressions (Tom)
  If the outer sub-select has no direct dependency on the parent query, but the inner one does, the outer value might not get recalculated for new parent query rows.

• Fix planner to estimate that GROUP BY expressions yielding boolean results always result in two groups, regardless of the expressions' contents (Tom)
  This is very substantially more accurate than the regular GROUP BY estimate for certain boolean tests like col IS NULL.

• Fix PL/pgSQL to not fail when a FOR loop's target variable is a record containing composite-type fields (Tom)

• Fix PL/Tcl to behave correctly with Tcl 8.5, and to be more careful about the encoding of data sent to or from Tcl (Tom)

• Improve performance of PQescapeBytea() (Rudolf Leitgeb)

• On Windows, work around a Microsoft bug by preventing libpq from trying to send more than 64kB per system call (Magnus)

• Fix ecpg to handle variables properly in SET commands (Michael)

• Improve pg_dump and pg_restore's error reporting after failure to send a SQL command (Tom)

• Fix pg_ctl to properly preserve postmaster command-line arguments across a restart (Bruce)

• Fix erroneous WAL file cutoff point calculation in pg_standby (Simon)

• Update time zone data files to tzdata release 2008f (for DST law changes in Argentina, Bahamas, Brazil, Mauritius, Morocco, Pakistan, Palestine, and Paraguay)

E.202. Release 8.3.3

Release date: 2008-06-12

This release contains one serious and one minor bug fix over 8.3.2. For information about new features in the 8.3 major release, see Section E.205.

E.202.1. Migration to Version 8.3.3

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.1, see Section E.204.
E.202.2. Changes

- Make `pg_get_ruledef()` parenthesize negative constants (Tom)

Before this fix, a negative constant in a view or rule might be dumped as, say, `-42::integer`, which is subtly incorrect: it should be `(-42)::integer` due to operator precedence rules. Usually this would make little difference, but it could interact with another recent patch to cause PostgreSQL to reject what had been a valid `SELECT DISTINCT` view query. Since this could result in `pg_dump` output failing to reload, it is being treated as a high-priority fix. The only released versions in which dump output is actually incorrect are 8.3.1 and 8.2.7.

- Make `ALTER AGGREGATE ... OWNER TO` update `pg_shdepend` (Tom)

This oversight could lead to problems if the aggregate was later involved in a `DROP OWNED` or `REASSIGN OWNED` operation.

E.203. Release 8.3.2

Release date: never released

This release contains a variety of fixes from 8.3.1. For information about new features in the 8.3 major release, see Section E.205.

E.203.1. Migration to Version 8.3.2

A dump/restore is not required for those running 8.3.X. However, if you are upgrading from a version earlier than 8.3.1, see Section E.204.

E.203.2. Changes

- Fix `ERRORDATA_STACK_SIZE exceeded` crash that occurred on Windows when using UTF-8 database encoding and a different client encoding (Tom)

- Fix incorrect archive truncation point calculation for the `%r` macro in `restore_command` parameters (Simon)

  This could lead to data loss if a warm-standby script relied on `%r` to decide when to throw away WAL segment files.

- Fix `ALTER TABLE ADD COLUMN ... PRIMARY KEY` so that the new column is correctly checked to see if it’s been initialized to all non-nulls (Brendan Jurd)

  Previous versions neglected to check this requirement at all.

- Fix `REASSIGN OWNED` so that it works on procedural languages too (Alvaro)

- Fix problems with `SELECT FOR UPDATE/SHARE occurring as a subquery in a query with a non-SELECT top-level operation` (Tom)

- Fix possible `CREATE TABLE` failure when inheriting the “same” constraint from multiple parent relations that inherited that constraint from a common ancestor (Tom)

- Fix `pg_get_ruledef()` to show the alias, if any, attached to the target table of an `UPDATE` or `DELETE` (Tom)

- Restore the pre-8.3 behavior that an out-of-range block number in a TID being used in a TidScan plan results in silently not matching any rows (Tom)
8.3.0 and 8.3.1 threw an error instead.

- Fix GIN bug that could result in a too many LWLocks taken failure (Teodor)
- Fix broken GiST comparison function for tsquery (Teodor)
- Fix tsvector_update_trigger() and ts_stat() to accept domains over the types they expect to work with (Tom)
- Fix failure to support enum data types as foreign keys (Tom)
- Avoid possible crash when decompressing corrupted data (Zdenek Kotala)
- Fix race conditions between delayed unlinks and DROP DATABASE (Heikki)

In the worst case this could result in deleting a newly created table in a new database that happened to get the same OID as the recently-dropped one; but of course that is an extremely low-probability scenario.

- Repair two places where SIGTERM exit of a backend could leave corrupted state in shared memory (Tom)

Neither case is very important if SIGTERM is used to shut down the whole database cluster together, but there was a problem if someone tried to SIGTERM individual backends.

- Fix possible crash due to incorrect plan generated for an \( x \text{ IN } \{ \text{SELECT } y \text{ FROM } \ldots \} \) clause when \( x \) and \( y \) have different data types; and make sure the behavior is semantically correct when the conversion from \( y \)'s type to \( x \)'s type is lossy (Tom)
- Fix oversight that prevented the planner from substituting known Param values as if they were constants (Tom)

This mistake partially disabled optimization of unnamed extended-Query statements in 8.3.0 and 8.3.1: in particular the LIKE-to-indexscan optimization would never be applied if the LIKE pattern was passed as a parameter, and constraint exclusion depending on a parameter value didn't work either.

- Fix planner failure when an indexable MIN or MAX aggregate is used with DISTINCT or ORDER BY (Tom)
- Fix planner to ensure it never uses a "physical tlist" for a plan node that is feeding a Sort node (Tom)

This led to the sort having to push around more data than it really needed to, since unused column values were included in the sorted data.

- Avoid unnecessary copying of query strings (Tom)

This fixes a performance problem introduced in 8.3.0 when a very large number of commands are submitted as a single query string.

- Make TransactionIdIsCurrentTransactionId() use binary search instead of linear search when checking child-transaction XIDs (Heikki)

This fixes some cases in which 8.3.0 was significantly slower than earlier releases.

- Fix conversions between ISO-8859-5 and other encodings to handle Cyrillic “Yo” characters (e and E with two dots) (Sergey Burladyan)
- Fix several datatype input functions, notably array_in(), that were allowing unused bytes in their results to contain uninitialized, unpredictable values (Tom)
This could lead to failures in which two apparently identical literal values were not seen as equal, resulting in the parser complaining about unmatched ORDER BY and DISTINCT expressions.

- Fix a corner case in regular-expression substring matching (substring(string from pattern)) (Tom)

The problem occurs when there is a match to the pattern overall but the user has specified a parenthesized subexpression and that subexpression hasn't got a match. An example is substring('foo from 'foo(bar)?'). This should return NULL, since (bar) isn't matched, but it was mistakenly returning the whole-pattern match instead (ie, foo).

- Prevent cancellation of an auto-vacuum that was launched to prevent XID wraparound (Alvaro)

- Improve ANALYZE's handling of in-doubt tuples (those inserted or deleted by a not-yet-committed transaction) so that the counts it reports to the stats collector are more likely to be correct (Pavan Deolasee)

- Fix initdb to reject a relative path for its --xlogdir (-X) option (Tom)

- Make psql print tab characters as an appropriate number of spaces, rather than \x09 as was done in 8.3.0 and 8.3.1 (Bruce)

- Update time zone data files to tzdata release 2008c (for DST law changes in Morocco, Iraq, Chibalsan, Pakistan, Syria, Cuba, and Argentina/San_Luis)

- Add ECPGget_PGconn() function to ecpglib (Michael)

- Fix incorrect result from ecpg's PGTYPEStimestamp_sub() function (Michael)

- Fix handling of continuation line markers in ecpg (Michael)

- Fix possible crashes in contrib/cube functions (Tom)

- Fix core dump in contrib/xml2's xpath_table() function when the input query returns a NULL value (Tom)

- Fix contrib/xml2's makefile to not override CFLAGS, and make it auto-configure properly for libxslt present or not (Tom)

**E.204. Release 8.3.1**

Release date: 2008-03-17

This release contains a variety of fixes from 8.3.0. For information about new features in the 8.3 major release, see Section E.205.

**E.204.1. Migration to Version 8.3.1**

A dump/restore is not required for those running 8.3.X. However, you might need to REINDEX indexes on textual columns after updating, if you are affected by the Windows locale issue described below.

**E.204.2. Changes**

- Fix character string comparison for Windows locales that consider different character combinations as equal (Tom)

This fix applies only on Windows and only when using UTF-8 database encoding. The same fix was made for all other cases over two years ago, but Windows with UTF-8 uses a separate code path
that was not updated. If you are using a locale that considers some non-identical strings as equal, you may need to \texttt{REINDEX} to fix existing indexes on textual columns.

- \textit{Repair corner-case bugs in \texttt{VACUUM FULL}} \texttt{(Tom)}

  A potential deadlock between concurrent \texttt{VACUUM FULL} operations on different system catalogs was introduced in 8.2. This has now been corrected. 8.3 made this worse because the deadlock could occur within a critical code section, making it a \texttt{PANIC} rather than just \texttt{ERROR} condition.

  Also, a \texttt{VACUUM FULL} that failed partway through vacuuming a system catalog could result in cache corruption in concurrent database sessions.

  Another \texttt{VACUUM FULL} bug introduced in 8.3 could result in a crash or out-of-memory report when dealing with pages containing no live tuples.

- \textit{Fix misbehavior of foreign key checks involving \texttt{character} or \texttt{bit} columns} \texttt{(Tom)}

  If the referencing column were of a different but compatible type (for instance \texttt{varchar}), the constraint was enforced incorrectly.

- \textit{Avoid needless deadlock failures in no-op foreign-key checks} \texttt{(Stephan Szabo, Tom)}

- \textit{Fix possible core dump when re-planning a prepared query} \texttt{(Tom)}

  This bug affected only protocol-level prepare operations, not \texttt{SQL PREPARE}, and so tended to be seen only with JDBC, DBI, and other client-side drivers that use prepared statements heavily.

- \textit{Fix possible failure when re-planning a query that calls an SPI-using function} \texttt{(Tom)}

- \textit{Fix failure in row-wise comparisons involving columns of different datatypes} \texttt{(Tom)}

- \textit{Fix longstanding \texttt{LISTEN}/\texttt{NOTIFY} race condition} \texttt{(Tom)}

  In rare cases a session that had just executed a \texttt{LISTEN} might not get a notification, even though one would be expected because the concurrent transaction executing \texttt{NOTIFY} was observed to commit later.

  A side effect of the fix is that a transaction that has executed a not-yet-committed \texttt{LISTEN} command will not see any row in \texttt{pg_listener} for the \texttt{LISTEN}, should it choose to look; formerly it would have. This behavior was never documented one way or the other, but it is possible that some applications depend on the old behavior.

- \textit{Disallow \texttt{LISTEN} and \texttt{UNLISTEN} within a prepared transaction} \texttt{(Tom)}

  This was formerly allowed but trying to do it had various unpleasant consequences, notably that the originating backend could not exit as long as an \texttt{UNLISTEN} remained uncommitted.

- \textit{Disallow dropping a temporary table within a prepared transaction} \texttt{(Heikki)}

  This was correctly disallowed by 8.1, but the check was inadvertently broken in 8.2 and 8.3.

- \textit{Fix rare crash when an error occurs during a query using a hash index} \texttt{(Heikki)}

- \textit{Fix incorrect comparison of tsquery values} \texttt{(Teodor)}

- \textit{Fix incorrect behavior of \texttt{LIKE} with non-ASCII characters in single-byte encodings} \texttt{(Rolf Jentsch)}

- \textit{Disable xmlvalidate} \texttt{(Tom)}

  This function should have been removed before 8.3 release, but was inadvertently left in the source code. It poses a small security risk since unprivileged users could use it to read the first few characters of any file accessible to the server.
• Fix memory leaks in certain usages of set-returning functions (Neil)

• Make `encode( bytea, 'escape')` convert all high-bit-set byte values into \nnn octal escape sequences (Tom)

  This is necessary to avoid encoding problems when the database encoding is multi-byte. This change could pose compatibility issues for applications that are expecting specific results from `encode`.

• Fix input of datetime values for February 29 in years BC (Tom)

  The former coding was mistaken about which years were leap years.

• Fix “unrecognized node type” error in some variants of `ALTER OWNER` (Tom)

• Avoid tablespace permissions errors in `CREATE TABLE LIKE INCLUDING INDEXES` (Tom)

• Ensure `pg_stat_activity.waiting` flag is cleared when a lock wait is aborted (Tom)

• Fix handling of process permissions on Windows Vista (Dave, Magnus)

  In particular, this fix allows starting the server as the Administrator user.

• Update time zone data files to `tzdata` release 2008a (in particular, recent Chile changes); adjust timezone abbreviation VET (Venezuela) to mean UTC-4:30, not UTC-4:00 (Tom)

• Fix `ecpg` problems with arrays (Michael)

• Fix `pg_ctl` to correctly extract the postmaster's port number from command-line options (Itagaki Takahiro, Tom)

  Previously, `pg_ctl start -w` could try to contact the postmaster on the wrong port, leading to bogus reports of startup failure.

• Use `-fwrapv` to defend against possible misoptimization in recent gcc versions (Tom)

  This is known to be necessary when building PostgreSQL with gcc 4.3 or later.

• Enable building `contrib/uuid-ossp` with MSVC (Hirosi Saito)

E.205. Release 8.3

Release date: 2008-02-04

E.205.1. Overview

With significant new functionality and performance enhancements, this release represents a major leap forward for PostgreSQL. This was made possible by a growing community that has dramatically accelerated the pace of development. This release adds the following major features:

• Full text search is integrated into the core database system

• Support for the SQL/XML standard, including new operators and an XML data type

• Enumerated data types (`ENUM`)

• Arrays of composite types

• Universally Unique Identifier (`UUID`) data type

• Add control over whether NULLs sort first or last
• Updatable cursors
• Server configuration parameters can now be set on a per-function basis
• User-defined types can now have type modifiers
• Automatically re-plan cached queries when table definitions change or statistics are updated
• Numerous improvements in logging and statistics collection
• Support Security Service Provider Interface (SSPI) for authentication on Windows
• Support multiple concurrent autovacuum processes, and other autovacuum improvements
• Allow the whole PostgreSQL distribution to be compiled with Microsoft Visual C++

Major performance improvements are listed below. Most of these enhancements are automatic and do not require user changes or tuning:
• Asynchronous commit delays writes to WAL during transaction commit
• Checkpoint writes can be spread over a longer time period to smooth the I/O spike during each checkpoint
• Heap-Only Tuples (HOT) accelerate space reuse for most UPDATEs and DELETEs
• Just-in-time background writer strategy improves disk write efficiency
• Using non-persistent transaction IDs for read-only transactions reduces overhead and VACUUM requirements
• Per-field and per-row storage overhead has been reduced
• Large sequential scans no longer force out frequently used cached pages
• Concurrent large sequential scans can now share disk reads
• ORDER BY ... LIMIT can be done without sorting

The above items are explained in more detail in the sections below.

E.205.2. Migration to Version 8.3

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

E.205.2.1. General

• Non-character data types are no longer automatically cast to TEXT (Peter, Tom)

Previously, if a non-character value was supplied to an operator or function that requires text input, it was automatically cast to text, for most (though not all) built-in data types. This no longer happens: an explicit cast to text is now required for all non-character-string types. For example, these expressions formerly worked:

```
substr(current_date, 1, 4)
23 LIKE '2%'
```

but will now draw “function does not exist” and “operator does not exist” errors respectively. Use an explicit cast instead:
(Of course, you can use the more verbose `CAST()` syntax too.) The reason for the change is that these automatic casts too often caused surprising behavior. An example is that in previous releases, this expression was accepted but did not do what was expected:

```
current_date < 2017-11-17
```

This is actually comparing a date to an integer, which should be (and now is) rejected — but in the presence of automatic casts both sides were cast to `text` and a textual comparison was done, because the `text < text` operator was able to match the expression when no other `<` operator could.

Types `char(n)` and `varchar(n)` still cast to `text` automatically. Also, automatic casting to `text` still works for inputs to the concatenation (`||`) operator, so long as least one input is a character-string type.

- Full text search features from `contrib/tsearch2` have been moved into the core server, with some minor syntax changes

  `contrib/tsearch2` now contains a compatibility interface.

- `ARRAY(SELECT ...),` where the `SELECT` returns no rows, now returns an empty array, rather than `NULL` (Tom)

- The array type name for a base data type is no longer always the base type's name with an underscore prefix

  The old naming convention is still honored when possible, but application code should no longer depend on it. Instead use the new `pg_type.typarray` column to identify the array data type associated with a given type.

- `ORDER BY ... USING operator` must now use a less-than or greater-than `operator` that is defined in a btree operator class

  This restriction was added to prevent inconsistent results.

- `SET LOCAL` changes now persist until the end of the outermost transaction, unless rolled back (Tom)

  Previously `SET LOCAL`'s effects were lost after subtransaction commit (`RELEASE SAVEPOINT` or exit from a PL/pgSQL exception block).

- Commands rejected in transaction blocks are now also rejected in multiple-statement query strings (Tom)

  For example, "BEGIN; DROP DATABASE; COMMIT" will now be rejected even if submitted as a single query message.

- `ROLLBACK` outside a transaction block now issues `NOTICE` instead of `WARNING` (Bruce)

- Prevent `NOTIFY/LISTEN/UNLISTEN` from accepting schema-qualified names (Bruce)

  Formerly, these commands accepted `schema.relation` but ignored the schema part, which was confusing.

- `ALTER SEQUENCE` no longer affects the sequence's `currval()` state (Tom)
• Foreign keys now must match indexable conditions for cross-data-type references (Tom)

This improves semantic consistency and helps avoid performance problems.

• Restrict object size functions to users who have reasonable permissions to view such information (Tom)

For example, `pg_database_size()` now requires CONNECT permission, which is granted to everyone by default. `pg_tablespace_size()` requires CREATE permission in the tablespace, or is allowed if the tablespace is the default tablespace for the database.

• Remove the undocumented `!=` (not in) operator (Tom)

`NOT IN (SELECT ...)` is the proper way to perform this operation.

• Internal hashing functions are now more uniformly-distributed (Tom)

If application code was computing and storing hash values using internal PostgreSQL hashing functions, the hash values must be regenerated.

• C-code conventions for handling variable-length data values have changed (Greg Stark, Tom)

The new `SET_VARSIZE()` macro must be used to set the length of generated `varlena` values. Also, it might be necessary to expand (“de-TOAST”) input values in more cases.

• Continuous archiving no longer reports each successful archive operation to the server logs unless `DEBUG` level is used (Simon)

E.205.2.2. Configuration Parameters

• Numerous changes in administrative server parameters

`bgwriter_lru_percent`, `bgwriter_all_percent`, `bgwriter_all_maxpages`, `stats_start_collector` and `stats_reset_on_server_start` are removed. `redirect_stderr` is renamed to `logging_collector`. `stats_command_string` is renamed to `track_activities`. `stats_block_level` and `stats_row_level` are merged into `track_counts`. A new boolean configuration parameter, `archive_mode`, controls archiving. Autovacuum’s default settings have changed.

• Remove `stats_start_collector` parameter (Tom)

We now always start the collector process, unless UDP socket creation fails.

• Remove `stats_reset_on_server_start` parameter (Tom)

This was removed because `pg_stat_reset()` can be used for this purpose.

• Commenting out a parameter in `postgresql.conf` now causes it to revert to its default value (Joachim Wieland)

Previously, commenting out an entry left the parameter’s value unchanged until the next server restart.

E.205.2.3. Character Encodings

• Add more checks for invalidly-encoded data (Andrew)

This change plugs some holes that existed in literal backslash escape string processing and COPY escape processing. Now the de-escaped string is rechecked to see if the result created an invalid multi-byte character.
• Disallow database encodings that are inconsistent with the server's locale setting (Tom)

On most platforms, C locale is the only locale that will work with any database encoding. Other locale settings imply a specific encoding and will misbehave if the database encoding is something different. (Typical symptoms include bogus textual sort order and wrong results from upper() or lower().) The server now rejects attempts to create databases that have an incompatible encoding.

• Ensure that chr() cannot create invalidly-encoded values (Andrew)

In UTF8-encoded databases the argument of chr() is now treated as a Unicode code point. In other multi-byte encodings chr()’s argument must designate a 7-bit ASCII character. Zero is no longer accepted. ascii() has been adjusted to match.

• Adjust convert() behavior to ensure encoding validity (Andrew)

The two argument form of convert() has been removed. The three argument form now takes a bytea first argument and returns a bytea. To cover the loss of functionality, three new functions have been added:

• convert_from(bytea, name) returns text — converts the first argument from the named encoding to the database encoding

• convert_to(text, name) returns bytea — converts the first argument from the database encoding to the named encoding

• length(bytea, name) returns integer — gives the length of the first argument in characters in the named encoding

• Remove convert(argument USING conversion_name) (Andrew)

Its behavior did not match the SQL standard.

• Make JOHAB encoding client-only (Tatsuo)

JOHAB is not safe as a server-side encoding.

E.205.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 8.3 and the previous major release.

E.205.3.1. Performance

• Asynchronous commit delays writes to WAL during transaction commit (Simon)

This feature dramatically increases performance for short data-modifying transactions. The disadvantage is that because disk writes are delayed, if the database or operating system crashes before data is written to the disk, committed data will be lost. This feature is useful for applications that can accept some data loss. Unlike turning off fsync, using asynchronous commit does not put database consistency at risk; the worst case is that after a crash the last few reportedly-committed transactions might not be committed after all. This feature is enabled by turning off synchronous_commit (which can be done per-session or per-transaction, if some transactions are critical and others are not). wal_writer_delay can be adjusted to control the maximum delay before transactions actually reach disk.

• Checkpoint writes can be spread over a longer time period to smooth the I/O spike during each checkpoint (Itagaki Takahiro and Heikki Linnakangas)

Previously all modified buffers were forced to disk as quickly as possible during a checkpoint, causing an I/O spike that decreased server performance. This new approach spreads out disk writes
during checkpoints, reducing peak I/O usage. (User-requested and shutdown checkpoints are still written as quickly as possible.)

- Heap-Only Tuples (HOT) accelerate space reuse for most UPDATEs and DELETEs (Pavan Deolasee, with ideas from many others)

  UPDATEs and DELETEs leave dead tuples behind, as do failed INSERTs. Previously only VACUUM could reclaim space taken by dead tuples. With HOT dead tuple space can be automatically reclaimed at the time of INSERT or UPDATE if no changes are made to indexed columns. This allows for more consistent performance. Also, HOT avoids adding duplicate index entries.

- Just-in-time background writer strategy improves disk write efficiency (Greg Smith, Itagaki Takahiro)

  This greatly reduces the need for manual tuning of the background writer.

- Per-field and per-row storage overhead have been reduced (Greg Stark, Heikki Linnakangas)

  Variable-length data types with data values less than 128 bytes long will see a storage decrease of 3 to 6 bytes. For example, two adjacent char(1) fields now use 4 bytes instead of 16. Row headers are also 4 bytes shorter than before.

- Using non-persistent transaction IDs for read-only transactions reduces overhead and VACUUM requirements (Florian Pflug)

  Non-persistent transaction IDs do not increment the global transaction counter. Therefore, they reduce the load on pg_clog and increase the time between forced vacuums to prevent transaction ID wraparound. Other performance improvements were also made that should improve concurrency.

- Avoid incrementing the command counter after a read-only command (Tom)

  There was formerly a hard limit of $2^{32}$ (4 billion) commands per transaction. Now only commands that actually changed the database count, so while this limit still exists, it should be significantly less annoying.

- Create a dedicated WAL writer process to off-load work from backends (Simon)

- Skip unnecessary WAL writes for CLUSTER and COPY (Simon)

  Unless WAL archiving is enabled, the system now avoids WAL writes for CLUSTER and just \texttt{fsync()}s the table at the end of the command. It also does the same for COPY if the table was created in the same transaction.

- Large sequential scans no longer force out frequently used cached pages (Simon, Heikki, Tom)

- Concurrent large sequential scans can now share disk reads (Jeff Davis)

  This is accomplished by starting the new sequential scan in the middle of the table (where another sequential scan is already in-progress) and wrapping around to the beginning to finish. This can affect the order of returned rows in a query that does not specify \texttt{ORDER BY}. The \texttt{syncrize_seqscans} configuration parameter can be used to disable this if necessary.

- \texttt{ORDER BY ... LIMIT} can be done without sorting (Greg Stark)

  This is done by sequentially scanning the table and tracking just the “top N” candidate rows, rather than performing a full sort of the entire table. This is useful when there is no matching index and the \texttt{LIMIT} is not large.

- Put a rate limit on messages sent to the statistics collector by backends (Tom)

  This reduces overhead for short transactions, but might sometimes increase the delay before statistics are tallied.
Release Notes

• Improve hash join performance for cases with many NULLs (Tom)
• Speed up operator lookup for cases with non-exact datatype matches (Tom)

E.205.3.2. Server

• Autovacuum is now enabled by default (Alvaro)

  Several changes were made to eliminate disadvantages of having autovacuum enabled, thereby justifying the change in default. Several other autovacuum parameter defaults were also modified.

• Support multiple concurrent autovacuum processes (Alvaro, Itagaki Takahiro)

  This allows multiple autovacuums to run concurrently. This prevents vacuuming of a large table from delaying vacuuming of smaller tables.

• Automatically re-plan cached queries when table definitions change or statistics are updated (Tom)

  Previously PL/pgSQL functions that referenced temporary tables would fail if the temporary table was dropped and recreated between function invocations, unless EXECUTE was used. This improvement fixes that problem and many related issues.

• Add a `temp_tablespaces` parameter to control the tablespaces for temporary tables and files (Jaime Casanova, Albert Cervera, Bernd Helmle)

  This parameter defines a list of tablespaces to be used. This enables spreading the I/O load across multiple tablespaces. A random tablespace is chosen each time a temporary object is created. Temporary files are no longer stored in per-database `pgsql_tmp/` directories but in per-tablespace directories.

• Place temporary tables’ TOAST tables in special schemas named `pg_toast_temp_nnn` (Tom)

  This allows low-level code to recognize these tables as temporary, which enables various optimizations such as not WAL-logging changes and using local rather than shared buffers for access. This also fixes a bug wherein backends unexpectedly held open file references to temporary TOAST tables.

• Fix problem that a constant flow of new connection requests could indefinitely delay the postmaster from completing a shutdown or a crash restart (Tom)

• Guard against a very-low-probability data loss scenario by preventing re-use of a deleted table’s relfilenode until after the next checkpoint (Heikki)

• Fix `CREATE CONSTRAINT TRIGGER` to convert old-style foreign key trigger definitions into regular foreign key constraints (Tom)

  This will ease porting of foreign key constraints carried forward from pre-7.3 databases, if they were never converted using `contrib/adddepend`.

• Fix `DEFAULT NULL` to override inherited defaults (Tom)

  `DEFAULT NULL` was formerly considered a noise phrase, but it should (and now does) override non-null defaults that would otherwise be inherited from a parent table or domain.

• Add new encodings EUC_JIS_2004 and SHIFT_JIS_2004 (Tatsuo)

  These new encodings can be converted to and from UTF-8.

• Change server startup log message from “database system is ready” to “database system is ready to accept connections”, and adjust its timing

  The message now appears only when the postmaster is really ready to accept connections.
E.205.3.3. Monitoring

- Add `log_autovacuum_min_duration` parameter to support configurable logging of auto-vacuum activity (Simon, Alvaro)
- Add `log_lock_waits` parameter to log lock waiting (Simon)
- Add `log_temp_files` parameter to log temporary file usage (Bill Moran)
- Add `log_checkpoints` parameter to improve logging of checkpoints (Greg Smith, Heikki)
- `log_line_prefix` now supports \%s and \%c escapes in all processes (Andrew)
  
  Previously these escapes worked only for user sessions, not for background database processes.
- Add `log_restartpoints` to control logging of point-in-time recovery restart points (Simon)
- Last transaction end time is now logged at end of recovery and at each logged restart point (Simon)
- Autovacuum now reports its activity start time in `pg_stat_activity` (Tom)
- Allow server log output in comma-separated value (CSV) format (Arul Shaji, Greg Smith, Andrew Dunstan)
  
  CSV-format log files can easily be loaded into a database table for subsequent analysis.
- Use PostgreSQL-supplied timezone support for formatting timestamps displayed in the server log (Tom)
  
  This avoids Windows-specific problems with localized time zone names that are in the wrong encoding. There is a new `log_timezone` parameter that controls the time zone used in log messages, independently of the client-visible `timezone` parameter.
- New system view `pg_stat_bgwriter` displays statistics about background writer activity (Magnus)
- Add new columns for database-wide tuple statistics to `pg_stat_database` (Magnus)
- Add an `xact_start` (transaction start time) column to `pg_stat_activity` (Neil)
  
  This makes it easier to identify long-running transactions.
- Add `n_live_tuples` and `n_dead_tuples` columns to `pg_stat_all_tables` and related views (Glen Parker)
- Merge `stats_block_level` and `stats_row_level` parameters into a single parameter `track_counts`, which controls all messages sent to the statistics collector process (Tom)
- Rename `stats_command_string` parameter to `track_activities` (Tom)
- Fix statistical counting of live and dead tuples to recognize that committed and aborted transactions have different effects (Tom)

E.205.3.4. Authentication

- Support Security Service Provider Interface (SSPI) for authentication on Windows (Magnus)
- Support GSSAPI authentication (Henry Hotz, Magnus)
  
  This should be preferred to native Kerberos authentication because GSSAPI is an industry standard.
• Support a global SSL configuration file (Victor Wagner)

• Add `ssl_ciphers` parameter to control accepted SSL ciphers (Victor Wagner)

• Add a Kerberos realm parameter, `krb_realm` (Magnus)

E.205.3.5. Write-Ahead Log (WAL) and Continuous Archiving

• Change the timestamps recorded in transaction WAL records from `time_t` to `TimestampTz` representation (Tom)

  This provides sub-second resolution in WAL, which can be useful for point-in-time recovery.

• Reduce WAL disk space needed by warm standby servers (Simon)

  This change allows a warm standby server to pass the name of the earliest still-needed WAL file to the recovery script, allowing automatic removal of no-longer-needed WAL files. This is done using `%r` in the `restore_command` parameter of `recovery.conf`.

• New boolean configuration parameter, `archive_mode`, controls archiving (Simon)

  Previously setting `archive_command` to an empty string turned off archiving. Now `archive_mode` turns archiving on and off, independently of `archive_command`. This is useful for stopping archiving temporarily.

E.205.3.6. Queries

• Full text search is integrated into the core database system (Teodor, Oleg)

  Text search has been improved, moved into the core code, and is now installed by default. `contrib/tsearch2` now contains a compatibility interface.

• Add control over whether NULLs sort first or last (Teodor, Tom)

  The syntax is `ORDER BY ... NULLS FIRST/LAST`.

• Allow per-column ascending/descending (`ASC/DESC`) ordering options for indexes (Teodor, Tom)

  Previously a query using `ORDER BY` with mixed `ASC/DESC` specifiers could not fully use an index. Now an index can be fully used in such cases if the index was created with matching `ASC/DESC` specifications. NULL sort order within an index can be controlled, too.

• Allow `col IS NULL` to use an index (Teodor)

• Updatable cursors (Arul Shaji, Tom)

  This eliminates the need to reference a primary key to `UPDATE` or `DELETE` rows returned by a cursor. The syntax is `UPDATE/DELETE WHERE CURRENT OF`.

• Allow `FOR UPDATE` in cursors (Arul Shaji, Tom)

• Create a general mechanism that supports casts to and from the standard string types (`TEXT`, `VARCHAR`, `CHAR`) for every datatype, by invoking the datatype's I/O functions (Tom)

  Previously, such casts were available only for types that had specialized function(s) for the purpose. These new casts are assignment-only in the to-string direction, explicit-only in the other direction, and therefore should create no surprising behavior.

• Allow `UNION` and related constructs to return a domain type, when all inputs are of that domain type (Tom)
Formerly, the output would be considered to be of the domain's base type.

- Allow limited hashing when using two different data types (Tom)

  This allows hash joins, hash indexes, hashed subplans, and hash aggregation to be used in situations involving cross-data-type comparisons, if the data types have compatible hash functions. Currently, cross-data-type hashing support exists for smallint/integer/bigint, and for float4/float8.

- Improve optimizer logic for detecting when variables are equal in a WHERE clause (Tom)

  This allows mergejoins to work with descending sort orders, and improves recognition of redundant sort columns.

- Improve performance when planning large inheritance trees in cases where most tables are excluded by constraints (Tom)

**E.205.3.7. Object Manipulation**

- Arrays of composite types (David Fetter, Andrew, Tom)

  In addition to arrays of explicitly-declared composite types, arrays of the rowtypes of regular tables and views are now supported, except for rowtypes of system catalogs, sequences, and TOAST tables.

- Server configuration parameters can now be set on a per-function basis (Tom)

  For example, functions can now set their own search_path to prevent unexpected behavior if a different search_path exists at run-time. Security definer functions should set search_path to avoid security loopholes.

- CREATE/ALTER FUNCTION now supports COST and ROWS options (Tom)

  COST allows specification of the cost of a function call. ROWS allows specification of the average number or rows returned by a set-returning function. These values are used by the optimizer in choosing the best plan.

- Implement CREATE TABLE LIKE ... INCLUDING INDEXES (Trevor Hardcastle, Nikhil Sontakke, Neil)

- Allow CREATE INDEX CONCURRENTLY to ignore transactions in other databases (Simon)

- Add ALTER VIEW ... RENAME TO and ALTER SEQUENCE ... RENAME TO (David Fetter, Neil)

  Previously this could only be done via ALTER TABLE ... RENAME TO.

- Make CREATE/DROP/RENAME DATABASE wait briefly for conflicting backends to exit before failing (Tom)

  This increases the likelihood that these commands will succeed.

- Allow triggers and rules to be deactivated in groups using a configuration parameter, for replication purposes (Jan)

  This allows replication systems to disable triggers and rewrite rules as a group without modifying the system catalogs directly. The behavior is controlled by ALTER TABLE and a new parameter session_replication_role.

- User-defined types can now have type modifiers (Teodor, Tom)
This allows a user-defined type to take a modifier, like \texttt{ssnum(7)}. Previously only built-in data types could have modifiers.

**E.205.3.8. Utility Commands**

- Non-superuser database owners now are able to add trusted procedural languages to their databases by default (Jeremy Drake)

  While this is reasonably safe, some administrators might wish to revoke the privilege. It is controlled by \texttt{pg\_pltemplate.tmpl\_db\_create}.

- Allow a session's current parameter setting to be used as the default for future sessions (Tom)

  This is done with \texttt{SET ... FROM CURRENT} in \texttt{CREATE/ALTER FUNCTION, ALTER DATABASE, or ALTER ROLE}.

- Implement new commands \texttt{DISCARD ALL, DISCARD PLANS, DISCARD TEMPORARY, CLOSE ALL, and DEALLOCATE ALL} (Marko Kreen, Neil)

  These commands simplify resetting a database session to its initial state, and are particularly useful for connection-pooling software.

- Make \texttt{CLUSTER} MVCC-safe (Heikki Linnakangas)

  Formerly, \texttt{CLUSTER} would discard all tuples that were committed dead, even if there were still transactions that should be able to see them under MVCC visibility rules.

- Add new \texttt{CLUSTER} syntax: \texttt{CLUSTER table USING index} (Holger Schurig)

  The old \texttt{CLUSTER} syntax is still supported, but the new form is considered more logical.

- Fix \texttt{EXPLAIN} so it can show complex plans more accurately (Tom)

  References to subplan outputs are now always shown correctly, instead of using \texttt{?columnN?} for complicated cases.

- Limit the amount of information reported when a user is dropped (Alvaro)

  Previously, dropping (or attempting to drop) a user who owned many objects could result in large \texttt{NOTICE} or \texttt{ERROR} messages listing all these objects; this caused problems for some client applications. The length of the message is now limited, although a full list is still sent to the server log.

**E.205.3.9. Data Types**

- Support for the SQL/XML standard, including new operators and an \texttt{XML} data type (Nikolay Samokhvalov, Pavel Stehule, Peter)

- Enumerated data types (\texttt{ENUM}) (Tom Dunstan)

  This feature provides convenient support for fields that have a small, fixed set of allowed values. An example of creating an \texttt{ENUM} type is \texttt{CREATE TYPE mood AS ENUM ('sad', 'ok', 'happy')}.

- Universally Unique Identifier (\texttt{UUID}) data type (Gevik Babakhani, Neil)

  This closely matches RFC 4122.

- Widen the \texttt{MONEY} data type to 64 bits (D'Arcy Cain)

  This greatly increases the range of supported \texttt{MONEY} values.
• Fix float4/float8 to handle Infinity and NaN (Not A Number) consistently (Bruce)
  The code formerly was not consistent about distinguishing Infinity from overflow conditions.
• Allow leading and trailing whitespace during input of boolean values (Neil)
• Prevent COPY from using digits and lowercase letters as delimiters (Tom)

E.205.3.10. Functions

• Add new regular expression functions regexp_matches(), regexp_split_to_array(),
  and regexp_split_to_table() (Jeremy Drake, Neil)
  These functions provide extraction of regular expression subexpressions and allow splitting a string
  using a POSIX regular expression.
• Add lo_truncate() for large object truncation (Kris Jurka)
• Implement width_bucket() for the float8 data type (Neil)
• Add pg_stat_clear_snapshot() to discard statistics snapshots collected during the current
  transaction (Tom)
  The first request for statistics in a transaction takes a statistics snapshot that does not change during
  the transaction. This function allows the snapshot to be discarded and a new snapshot loaded during
  the next statistics query. This is particularly useful for PL/pgSQL functions, which are confined to
  a single transaction.
• Add isodow option to EXTRACT() and date_part() (Bruce)
  This returns the day of the week, with Sunday as seven. (dow returns Sunday as zero.)
• Add ID (ISO day of week) and IDDD (ISO day of year) format codes for to_char(),
  to_date(), and to_timestamp() (Brendan Jurd)
• Make to_timestamp() and to_date() assume TM (trim) option for potentially vari-
  able-width fields (Bruce)
  This matches Oracle's behavior.
• Fix off-by-one conversion error in to_date()/to_timestamp() D (non-ISO day of week)
  fields (Bruce)
• Make randseed() return void, rather than a useless integer value (Neil)
• Add a hash function for NUMERIC (Neil)
  This allows hash indexes and hash-based plans to be used with NUMERIC columns.
• Improve efficiency of LIKE/ILIKE, especially for multi-byte character sets like UTF-8 (Andrew,
  Itagaki Takahiro)
• Make currtid() functions require SELECT privileges on the target table (Tom)
• Add several txid_*() functions to query active transaction IDs (Jan)
  This is useful for various replication solutions.

E.205.3.11. PL/pgSQL Server-Side Language

• Add scrollable cursor support, including directional control in FETCH (Pavel Stehule)
• Allow IN as an alternative to FROM in PL/pgSQL's FETCH statement, for consistency with the backend's FETCH command (Pavel Stehule)

• Add MOVE to PL/pgSQL (Magnus, Pavel Stehule, Neil)

• Implement RETURN QUERY (Pavel Stehule, Neil)

  This adds convenient syntax for PL/pgSQL set-returning functions that want to return the result of a query. RETURN QUERY is easier and more efficient than a loop around RETURN NEXT.

• Allow function parameter names to be qualified with the function's name (Tom)

  For example, myfunc.myvar. This is particularly useful for specifying variables in a query where the variable name might match a column name.

• Make qualification of variables with block labels work properly (Tom)

  Formerly, outer-level block labels could unexpectedly interfere with recognition of inner-level record or row references.

• Tighten requirements for FOR loop STEP values (Tom)

  Prevent non-positive STEP values, and handle loop overflows.

• Improve accuracy when reporting syntax error locations (Tom)

  E.205.3.12. Other Server-Side Languages

  • Allow type-name arguments to PL/Perl spi_prepare() to be data type aliases in addition to names found in pg_type (Andrew)

  • Allow type-name arguments to PL/Python plpy.prepare() to be data type aliases in addition to names found in pg_type (Andrew)

  • Allow type-name arguments to PL/Tcl spi_prepare to be data type aliases in addition to names found in pg_type (Andrew)

  • Enable PL/PythonU to compile on Python 2.5 (Marko Kreen)

  • Support a true PL/Python boolean type in compatible Python versions (Python 2.3 and later) (Marko Kreen)

  • Fix PL/Tcl problems with thread-enabled libtcl spawning multiple threads within the backend (Steve Marshall, Paul Bayer, Doug Knight)

    This caused all sorts of unpleasantness.

  E.205.3.13. psql

  • List disabled triggers separately in \d output (Brendan Jurd)

  • In \d patterns, always match $ literally (Tom)

  • Show aggregate return types in \da output (Greg Sabino Mullane)

  • Add the function's volatility status to the output of \df+ (Neil)

  • Add \prompt capability (Chad Wagner)

  • Allow \pset, \t, and \x to specify on or off, rather than just toggling (Chad Wagner)
Release Notes

- Add `\sleep` capability (Jan)
- Enable `\timing` output for `\copy` (Andrew)
- Improve `\timing` resolution on Windows (Itagaki Takahiro)
- Flush `\o` output after each backslash command (Tom)
- Correctly detect and report errors while reading a `-f` input file (Peter)
- Remove `-u` option (this option has long been deprecated) (Tom)

E.205.3.14. pg_dump

- Add `--tablespaces-only` and `--roles-only` options to `pg_dumpall` (Dave Page)
- Add an output file option to `pg_dumpall` (Dave Page)
  
  This is primarily useful on Windows, where output redirection of child `pg_dump` processes does not work.

- Allow `pg_dumpall` to accept an initial-connection database name rather than the default template (Dave Page)
- In `-n` and `-t` switches, always match `$` literally (Tom)
- Improve performance when a database has thousands of objects (Tom)
- Remove `-u` option (this option has long been deprecated) (Tom)

E.205.3.15. Other Client Applications

- In `initdb`, allow the location of the `pg_xlog` directory to be specified (Euler Taveira de Oliveira)
- Enable server core dump generation in `pg_regress` on supported operating systems (Andrew)
- Add a `-t` (timeout) parameter to `pg_ctl` (Bruce)
  
  This controls how long `pg_ctl` will wait when waiting for server startup or shutdown. Formerly the timeout was hard-wired as 60 seconds.

- Add a `pg_ctl` option to control generation of server core dumps (Andrew)
- Allow Control-C to cancel `clusterdb`, `reindexdb`, and `vacuumdb` (Itagaki Takahiro, Magnus)
- Suppress command tag output for `createdb`, `createuser`, `dropdb`, and `dropuser` (Peter)

  The `--quiet` option is ignored and will be removed in 8.4. Progress messages when acting on all databases now go to `stdout` instead of `stderr` because they are not actually errors.

E.205.3.16. libpq

- Interpret the `dbName` parameter of `PQsetdbLogin()` as a `conninfo` string if it contains an equals sign (Andrew)
  
  This allows use of `conninfo` strings in client programs that still use `PQsetdbLogin()`.

- Support a global SSL configuration file (Victor Wagner)
- Add environment variable `PGSSLKEY` to control SSL hardware keys (Victor Wagner)
• Add `lo_truncate()` for large object truncation (Kris Jurka)

• Add `PQconnectionNeedsPassword()` that returns true if the server required a password but none was supplied (Joe Conway, Tom)

  If this returns true after a failed connection attempt, a client application should prompt the user for a password. In the past applications have had to check for a specific error message string to decide whether a password is needed; that approach is now deprecated.

• Add `PQconnectionUsedPassword()` that returns true if the supplied password was actually used (Joe Conway, Tom)

  This is useful in some security contexts where it is important to know whether a user-supplied password is actually valid.

**E.205.3.17. ecpg**

• Use V3 frontend/backend protocol (Michael)

  This adds support for server-side prepared statements.

• Use native threads, instead of pthreads, on Windows (Magnus)

• Improve thread-safety of ecpglib (Itagaki Takahiro)

• Make the ecpg libraries export only necessary API symbols (Michael)

**E.205.3.18. Windows Port**

• Allow the whole PostgreSQL distribution to be compiled with Microsoft Visual C++ (Magnus and others)

  This allows Windows-based developers to use familiar development and debugging tools. Windows executables made with Visual C++ might also have better stability and performance than those made with other tool sets. The client-only Visual C++ build scripts have been removed.

• Drastically reduce postmaster's memory usage when it has many child processes (Magnus)

• Allow regression tests to be started by an administrative user (Magnus)

• Add native shared memory implementation (Magnus)

**E.205.3.19. Server Programming Interface (SPI)**

• Add cursor-related functionality in SPI (Pavel Stehule)

  Allow access to the cursor-related planning options, and add FETCH/MOVE routines.

• Allow execution of cursor commands through `SPI_execute` (Tom)

  The macro `SPI_ERROR_CURSOR` still exists but will never be returned.

• SPI plan pointers are now declared as `SPIPlanPtr` instead of `void *` (Tom)

  This does not break application code, but switching is recommended to help catch simple programming mistakes.

**E.205.3.20. Build Options**

• Add configure option `--enable-profiling` to enable code profiling (works only with gcc) (Korry Douglas and Nikhil Sontakke)
• Add configure option \--with-system-tzdata to use the operating system's time zone database (Peter)

• Fix PGXS so extensions can be built against PostgreSQL installations whose pg_config program does not appear first in the PATH (Tom)

• Support \texttt{gmake draft} when building the SGML documentation (Bruce)

Unless \texttt{draft} is used, the documentation build will now be repeated if necessary to ensure the index is up-to-date.

\textbf{E.205.3.21. Source Code}

• Rename macro \texttt{DLLIMPORT} to \texttt{PGDLLIMPORT} to avoid conflicting with third party includes (like Tcl) that define \texttt{DLLIMPORT} (Magnus)

• Create "operator families" to improve planning of queries involving cross-data-type comparisons (Tom)

• Update GIN \texttt{extractQuery()} API to allow signalling that nothing can satisfy the query (Teodor)

• Move \texttt{NAMEDATALEN} definition from \texttt{postgres_ext.h} to \texttt{pg_config_manual.h} (Peter)

• Provide \texttt{strlcpy()} and \texttt{strlcat()} on all platforms, and replace error-prone uses of \texttt{strncpy()}, \texttt{strncat()}, etc (Peter)

• Create hooks to let an external plugin monitor (or even replace) the planner and create plans for hypothetical situations (Gurjeet Singh, Tom)

• Create a function variable \texttt{join_search_hook} to let plugins override the join search order portion of the planner (Julius Stroffek)

• Add \texttt{tas()} support for Renesas' M32R processor (Kazuhiro Inaoka)

• \texttt{quote_identifier()} and \texttt{pg_dump} no longer quote keywords that are unreserved according to the grammar (Tom)

• Change the on-disk representation of the \texttt{NUMERIC} data type so that the \texttt{sign_dscale} word comes before the weight (Tom)

• Use SYSV semaphores rather than POSIX on Darwin \texttt{>= 6.0}, i.e., macOS 10.2 and up (Chris Marcellino)

• Add acronym and NFS documentation sections (Bruce)

• "Postgres" is now documented as an accepted alias for "PostgreSQL" (Peter)

• Add documentation about preventing database server spoofing when the server is down (Bruce)

\textbf{E.205.3.22. Contrib}

• Move \texttt{contrib} README content into the main PostgreSQL documentation (Albert Cervera i Areny)

• Add \texttt{contrib/pageinspect} module for low-level page inspection (Simon, Heikki)

• Add \texttt{contrib/pg_standby} module for controlling warm standby operation (Simon)

• Add \texttt{contrib/uuid-oss} module for generating UUID values using the OSSP UUID library (Peter)
Use configure --with-ossp-uuid to activate. This takes advantage of the new UUID builtin type.

- Add contrib/dict_int, contrib/dict_xsyn, and contrib/test_parser modules to provide sample add-on text search dictionary templates and parsers (Sergey Karpov)
- Allow contrib/pgbench to set the fillfactor (Pavan Deolasee)
- Add timestamps to contrib/pgbench -l (Greg Smith)
- Add usage count statistics to contrib/pgbuffercache (Greg Smith)
- Add GIN support for contrib/hstore (Teodor)
- Add GIN support for contrib/pg_trgm (Guillaume Smet, Teodor)
- Update OS/X startup scripts in contrib/start-scripts (Mark Cotner, David Fetter)
- Restrict pgrowlocks() and dblink_get_pkey() to users who have SELECT privilege on the target table (Tom)
- Restrict contrib/pgstattuple functions to superusers (Tom)
- contrib/xml2 is deprecated and planned for removal in 8.4 (Peter)

The new XML support in core PostgreSQL supersedes this module.

### E.206. Release 8.2.23

**Release date:** 2011-12-05

This release contains a variety of fixes from 8.2.22. For information about new features in the 8.2 major release, see Section E.229.

This is expected to be the last PostgreSQL release in the 8.2.X series. Users are encouraged to update to a newer release branch soon.

#### E.206.1. Migration to Version 8.2.23

A dump/restore is not required for those running 8.2.X.

However, a longstanding error was discovered in the definition of the information_schema.referential_constraints view. If you rely on correct results from that view, you should replace its definition as explained in the first changelog item below.

Also, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

#### E.206.2. Changes

- Fix bugs in information_schema.referential_constraints view (Tom Lane)

This view was being insufficiently careful about matching the foreign-key constraint to the depended-on primary or unique key constraint. That could result in failure to show a foreign key constraint at all, or showing it multiple times, or claiming that it depends on a different constraint than the one it really does.

Since the view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can (as a superuser) drop the information_schema schema then re-create it by sourcing $SHAREDIR/informa-
• Fix TOAST-related data corruption during CREATE TABLE dest AS SELECT * FROM src or INSERT INTO dest SELECT * FROM src (Tom Lane)

If a table has been modified by ALTER TABLE ADD COLUMN, attempts to copy its data verbatim to another table could produce corrupt results in certain corner cases. The problem can only manifest in this precise form in 8.4 and later, but we patched earlier versions as well in case there are other code paths that could trigger the same bug.

• Fix race condition during toast table access from stale syscache entries (Tom Lane)

The typical symptom was transient errors like “missing chunk number 0 for toast value NNNNN in pg_toast_2619”, where the cited toast table would always belong to a system catalog.

• Improve locale support in money type’s input and output (Tom Lane)

Aside from not supporting all standard lc_monetary formatting options, the input and output functions were inconsistent, meaning there were locales in which dumped money values could not be re-read.

• Don’t let transform_null_equals affect CASE foo WHEN NULL ... constructs (Heikki Linnakangas)

transform_null_equals is only supposed to affect foo = NULL expressions written directly by the user, not equality checks generated internally by this form of CASE.

• Change foreign-key trigger creation order to better support self-referential foreign keys (Tom Lane)

For a cascading foreign key that references its own table, a row update will fire both the ON UPDATE trigger and the CHECK trigger as one event. The ON UPDATE trigger must execute first, else the CHECK will check a non-final state of the row and possibly throw an inappropriate error. However, the firing order of these triggers is determined by their names, which generally sort in creation order since the triggers have auto-generated names following the convention “RI_ConstraintTrigger_NNNN”. A proper fix would require modifying that convention, which we will do in 9.2, but it seems risky to change it in existing releases. So this patch just changes the creation order of the triggers. Users encountering this type of error should drop and re-create the foreign key constraint to get its triggers into the right order.

• Preserve blank lines within commands in psql’s command history (Robert Haas)

The former behavior could cause problems if an empty line was removed from within a string literal, for example.

• Use the preferred version of xsubpp to build PL/Perl, not necessarily the operating system’s main copy (David Wheeler and Alex Hunsaker)

• Honor query cancel interrupts promptly in pgstatindex() (Robert Haas)

• Ensure VPATH builds properly install all server header files (Peter Eisentraut)

• Shorten file names reported in verbose error messages (Peter Eisentraut)

Regular builds have always reported just the name of the C file containing the error message call, but VPATH builds formerly reported an absolute path name.

• Fix interpretation of Windows timezone names for Central America (Tom Lane)

Map “Central America Standard Time” to CST6, not CST6CDT, because DST is generally not observed anywhere in Central America.
• Update time zone data files to tzdata release 2011n for DST law changes in Brazil, Cuba, Fiji, Palestine, Russia, and Samoa; also historical corrections for Alaska and British East Africa.

E.207. Release 8.2.22

Release date: 2011-09-26

This release contains a variety of fixes from 8.2.21. For information about new features in the 8.2 major release, see Section E.229.

The PostgreSQL community will stop releasing updates for the 8.2.X release series in December 2011. Users are encouraged to update to a newer release branch soon.

E.207.1. Migration to Version 8.2.22

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.207.2. Changes

• Fix multiple bugs in GiST index page split processing (Heikki Linnakangas)

The probability of occurrence was low, but these could lead to index corruption.

• Avoid possibly accessing off the end of memory in `ANALYZE` (Noah Misch)

This fixes a very-low-probability server crash scenario.

• Fix race condition in relcache init file invalidation (Tom Lane)

There was a window wherein a new backend process could read a stale init file but miss the inval messages that would tell it the data is stale. The result would be bizarre failures in catalog accesses, typically “could not read block 0 in file ...” later during startup.

• Fix memory leak at end of a GiST index scan (Tom Lane)

Commands that perform many separate GiST index scans, such as verification of a new GiST-based exclusion constraint on a table already containing many rows, could transiently require large amounts of memory due to this leak.

• Fix performance problem when constructing a large, lossy bitmap (Tom Lane)

• Fix array- and path-creating functions to ensure padding bytes are zeroes (Tom Lane)

This avoids some situations where the planner will think that semantically-equal constants are not equal, resulting in poor optimization.

• Work around gcc 4.6.0 bug that breaks WAL replay (Tom Lane)

This could lead to loss of committed transactions after a server crash.

• Fix dump bug for `VALUES` in a view (Tom Lane)

• Disallow `SELECT FOR UPDATE/SHARE` on sequences (Tom Lane)

This operation doesn’t work as expected and can lead to failures.

• Defend against integer overflow when computing size of a hash table (Tom Lane)

• Fix portability bugs in use of credentials control messages for “peer” authentication (Tom Lane)
• Fix typo in `pg_srand48` seed initialization (Andres Freund)

This led to failure to use all bits of the provided seed. This function is not used on most platforms (only those without `srandom`), and the potential security exposure from a less-random-than-expected seed seems minimal in any case.

• Avoid integer overflow when the sum of `LIMIT` and `OFFSET` values exceeds $2^{63}$ (Heikki Linnakangas)

• Add overflow checks to `int4` and `int8` versions of `generate_series()` (Robert Haas)

• Fix trailing-zero removal in `to_char()` (Marti Raudsepp)

In a format with `FM` and no digit positions after the decimal point, zeroes to the left of the decimal point could be removed incorrectly.

• Fix `pg_size_pretty()` to avoid overflow for inputs close to $2^{63}$ (Tom Lane)

• Fix psql's counting of script file line numbers during `COPY` from a different file (Tom Lane)

• Fix `pg_restore`'s direct-to-database mode for `standard_conforming_strings` (Tom Lane)

`pg_restore` could emit incorrect commands when restoring directly to a database server from an archive file that had been made with `standard_conforming_strings` set to on.

• Fix write-past-buffer-end and memory leak in libpq's LDAP service lookup code (Albe Laurenz)

• In libpq, avoid failures when using nonblocking I/O and an SSL connection (Martin Pihlak, Tom Lane)

• Improve libpq's handling of failures during connection startup (Tom Lane)

In particular, the response to a server report of `fork()` failure during SSL connection startup is now saner.

• Make ecpglib write `double` values with 15 digits precision (Akira Kurosawa)

• Apply upstream fix for blowfish signed-character bug (CVE-2011-2483) (Tom Lane)

`contrib/pg_crypto`'s blowfish encryption code could give wrong results on platforms where `char` is signed (which is most), leading to encrypted passwords being weaker than they should be.

• Fix memory leak in `contrib/seg` (Heikki Linnakangas)

• Fix `pgstatindex()` to give consistent results for empty indexes (Tom Lane)

• Allow building with perl 5.14 (Alex Hunsaker)

• Update configure script's method for probing existence of system functions (Tom Lane)

The version of autoconf we used in 8.3 and 8.2 could be fooled by compilers that perform link-time optimization.

• Fix assorted issues with build and install file paths containing spaces (Tom Lane)

• Update time zone data files to tzdata release 2011i for DST law changes in Canada, Egypt, Russia, Samoa, and South Sudan.

E.208. Release 8.2.21

Release date: 2011-04-18
This release contains a variety of fixes from 8.2.20. For information about new features in the 8.2 major release, see Section E.229.

E.208.1. Migration to Version 8.2.21

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.208.2. Changes

• Avoid potential deadlock during catalog cache initialization (Nikhil Sontakke)
  
  In some cases the cache loading code would acquire share lock on a system index before locking the index's catalog. This could deadlock against processes trying to acquire exclusive locks in the other, more standard order.

• Fix dangling-pointer problem in \texttt{BEFORE ROW UPDATE} trigger handling when there was a concurrent update to the target tuple (Tom Lane)

  This bug has been observed to result in intermittent “cannot extract system attribute from virtual tuple” failures while trying to do \texttt{UPDATE RETURNING ctid}. There is a very small probability of more serious errors, such as generating incorrect index entries for the updated tuple.

• Disallow \texttt{DROP TABLE} when there are pending deferred trigger events for the table (Tom Lane)

  Formerly the \texttt{DROP} would go through, leading to “could not open relation with OID nnn” errors when the triggers were eventually fired.

• Fix PL/Python memory leak involving array slices (Daniel Popowich)

• Fix pg_restore to cope with long lines (over 1KB) in TOC files (Tom Lane)

• Put in more safeguards against crashing due to division-by-zero with overly enthusiastic compiler optimization (Aurelien Jarno)

• Support use of \texttt{dlopen()} in FreeBSD and OpenBSD on MIPS (Tom Lane)

  There was a hard-wired assumption that this system function was not available on MIPS hardware on these systems. Use a compile-time test instead, since more recent versions have it.

• Fix compilation failures on HP-UX (Heikki Linnakangas)

• Fix path separator used by pg_regress on Cygwin (Andrew Dunstan)

• Update time zone data files to \texttt{tzdata} release 2011f for DST law changes in Chile, Cuba, Falkland Islands, Morocco, Samoa, and Turkey; also historical corrections for South Australia, Alaska, and Hawaii.

E.209. Release 8.2.20

Release date: 2011-01-31

This release contains a variety of fixes from 8.2.19. For information about new features in the 8.2 major release, see Section E.229.

E.209.1. Migration to Version 8.2.20

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.
E.209.2. Changes

• Avoid failures when EXPLAIN tries to display a simple-form CASE expression (Tom Lane)

If the CASE's test expression was a constant, the planner could simplify the CASE into a form that confused the expression-display code, resulting in “unexpected CASE WHEN clause” errors.

• Fix assignment to an array slice that is before the existing range of subscripts (Tom Lane)

If there was a gap between the newly added subscripts and the first pre-existing subscript, the code miscalculated how many entries needed to be copied from the old array's null bitmap, potentially leading to data corruption or crash.

• Avoid unexpected conversion overflow in planner for very distant date values (Tom Lane)

The date type supports a wider range of dates than can be represented by the timestamp types, but the planner assumed it could always convert a date to timestamp with impunity.

• Fix pg_restore's text output for large objects (BLOBs) when standard_conforming_strings is on (Tom Lane)

Although restoring directly to a database worked correctly, string escaping was incorrect if pg_restore was asked for SQL text output and standard_conforming_strings had been enabled in the source database.

• Fix erroneous parsing of tsquery values containing ... & !(subexpression) | ... (Tom Lane)

Queries containing this combination of operators were not executed correctly. The same error existed in contrib/intarray's query_int type and contrib/ltree's ltxtquery type.

• Fix buffer overrun in contrib/intarray's input function for the query_int type (Apple)

This bug is a security risk since the function's return address could be overwritten. Thanks to Apple Inc's security team for reporting this issue and supplying the fix. (CVE-2010-4015)

• Fix bug in contrib/seg's GiST picksplit algorithm (Alexander Korotkov)

This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a seg column. If you have such an index, consider REINDEXing it after installing this update. (This is identical to the bug that was fixed in contrib/cube in the previous update.)

E.210. Release 8.2.19

Release date: 2010-12-16

This release contains a variety of fixes from 8.2.18. For information about new features in the 8.2 major release, see Section E.229.

E.210.1. Migration to Version 8.2.19

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.210.2. Changes

• Force the default wal_sync_method to be fdatasync on Linux (Tom Lane, Marti Raudsepp)
The default on Linux has actually been `fdatasync` for many years, but recent kernel changes caused PostgreSQL to choose `open_datasync` instead. This choice did not result in any performance improvement, and caused outright failures on certain filesystems, notably `ext4` with the `data=journal` mount option.

- Fix assorted bugs in WAL replay logic for GIN indexes (Tom Lane)
  This could result in “bad buffer id: 0” failures or corruption of index contents during replication.

- Fix recovery from base backup when the starting checkpoint WAL record is not in the same WAL segment as its redo point (Jeff Davis)

- Add support for detecting register-stack overrun on IA64 (Tom Lane)
  The IA64 architecture has two hardware stacks. Full prevention of stack-overrun failures requires checking both.

- Add a check for stack overflow in `copyObject()` (Tom Lane)
  Certain code paths could crash due to stack overflow given a sufficiently complex query.

- Fix detection of page splits in temporary GiST indexes (Heikki Linnakangas)
  It is possible to have a “concurrent” page split in a temporary index, if for example there is an open cursor scanning the index when an insertion is done. GiST failed to detect this case and hence could deliver wrong results when execution of the cursor continued.

- Avoid memory leakage while ANALYZE'ing complex index expressions (Tom Lane)

- Ensure an index that uses a whole-row Var still depends on its table (Tom Lane)
  An index declared like `create index i on t (foo(t.*))` would not automatically get dropped when its table was dropped.

- Do not “inline” a SQL function with multiple `OUT` parameters (Tom Lane)
  This avoids a possible crash due to loss of information about the expected result rowtype.

- Behave correctly if `ORDER BY`, `LIMIT`, `FOR UPDATE`, or `WITH` is attached to the VALUES part of `INSERT ... VALUES` (Tom Lane)

- Fix constant-folding of `COALESCE()` expressions (Tom Lane)
  The planner would sometimes attempt to evaluate sub-expressions that in fact could never be reached, possibly leading to unexpected errors.

- Add print functionality for `InhRelation` nodes (Tom Lane)
  This avoids a failure when `debug_print_parse` is enabled and certain types of query are executed.

- Fix incorrect calculation of distance from a point to a horizontal line segment (Tom Lane)
  This bug affected several different geometric distance-measurement operators.

- Fix PL/pgSQL's handling of “simple” expressions to not fail in recursion or error-recovery cases (Tom Lane)

- Fix PL/Python's handling of set-returning functions (Jan Urbanski)
  Attempts to call SPI functions within the iterator generating a set result would fail.
• Fix bug in contrib/cube's GiST picksplit algorithm (Alexander Korotkov)

This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a cube column. If you have such an index, consider REINDEXing it after installing this update.

• Don't emit “identifier will be truncated” notices in contrib/dblink except when creating new connections (Itagaki Takahiro)

• Fix potential coredump on missing public key in contrib/pgcrypto (Marti Raudsepp)

• Fix memory leak in contrib/xml2's XPath query functions (Tom Lane)

• Update time zone data files to tzdata release 2010e for DST law changes in Fiji and Samoa; also historical corrections for Hong Kong.

E.211. Release 8.2.18

Release date: 2010-10-04

This release contains a variety of fixes from 8.2.17. For information about new features in the 8.2 major release, see Section E.229.

E.211.1. Migration to Version 8.2.18

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.211.2. Changes

• Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function's owner.

The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

• Prevent possible crashes in pg_get_expr() by disallowing it from being called with an argument that is not one of the system catalog columns it's intended to be used with (Heikki Linnakangas, Tom Lane)

• Fix Windows shared-memory allocation code (Tsutomu Yamada, Magnus Hagander)

This bug led to the often-reported “could not reattach to shared memory” error message. This is a back-patch of a fix that was applied to newer branches some time ago.

• Treat exit code 128 (ERROR_WAIT_NO_CHILDREN) as non-fatal on Windows (Magnus Hagander)
Under high load, Windows processes will sometimes fail at startup with this error code. Formerly the postmaster treated this as a panic condition and restarted the whole database, but that seems to be an overreaction.

- Fix possible duplicate scans of UNION ALL member relations (Tom Lane)
- Fix “cannot handle unplanned sub-select” error (Tom Lane)
  This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.
- Reduce PANIC to ERROR in some occasionally-reported btree failure cases, and provide additional detail in the resulting error messages (Tom Lane)
  This should improve the system’s robustness with corrupted indexes.
- Prevent show_session_authorization() from crashing within autovacuum processes (Tom Lane)
- Defend against functions returning setof record where not all the returned rows are actually of the same rowtype (Tom Lane)
- Fix possible failure when hashing a pass-by-reference function result (Tao Ma, Tom Lane)
- Take care to fsync the contents of lockfiles (both postmaster.pid and the socket lockfile) while writing them (Tom Lane)
  This omission could result in corrupted lockfile contents if the machine crashes shortly after postmaster start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.
- Avoid recursion while assigning XIDs to heavily-nested subtransactions (Andres Freund, Robert Haas)
  The original coding could result in a crash if there was limited stack space.
- Fix log_line_prefix’s %i escape, which could produce junk early in backend startup (Tom Lane)
- Fix possible data corruption in ALTER TABLE ... SET TABLESPACE when archiving is enabled (Jeff Davis)
- Allow CREATE DATABASE and ALTER DATABASE ... SET TABLESPACE to be interrupted by query-cancel (Guillaume Lelarge)
- In PL/Python, defend against null pointer results from PyObject_AsVoidPtr and PyObject_FromVoidPtr (Peter Eisentraut)
- Improve contrib/dblink’s handling of tables containing dropped columns (Tom Lane)
- Fix connection leak after “duplicate connection name” errors in contrib/dblink (Itagaki Takahiro)
- Fix contrib/dblink to handle connection names longer than 62 bytes correctly (Itagaki Takahiro)
- Add hstore(text, text) function to contrib/hstore (Robert Haas)
  This function is the recommended substitute for the now-deprecated => operator. It was back-patched so that future-proofed code can be used with older server versions. Note that the patch will be effective only after contrib/hstore is installed or reinstalled in a particular database. Users might prefer to execute the CREATE FUNCTION command by hand, instead.
• Update build infrastructure and documentation to reflect the source code repository’s move from CVS to Git (Magnus Hagander and others)

• Update time zone data files to tzdata release 2010l for DST law changes in Egypt and Palestine; also historical corrections for Finland.

This change also adds new names for two Micronesian timezones: Pacific/Chuuk is now preferred over Pacific/Truk (and the preferred abbreviation is CHUT not TRUT) and Pacific/Pohnpei is preferred over Pacific/Ponape.

• Make Windows’ “N. Central Asia Standard Time” timezone map to Asia/Novosibirsk, not Asia/Almaty (Magnus Hagander)

Microsoft changed the DST behavior of this zone in the timezone update from KB976098. Asia/Novosibirsk is a better match to its new behavior.

E.212. Release 8.2.17

Release date: 2010-05-17

This release contains a variety of fixes from 8.2.16. For information about new features in the 8.2 major release, see Section E.229.

E.212.1. Migration to Version 8.2.17

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.212.2. Changes

• Enforce restrictions in plperl using an opmask applied to the whole interpreter, instead of using Safe.pm (Tim Bunce, Andrew Dunstan)

Recent developments have convinced us that Safe.pm is too insecure to rely on for making plperl trustable. This change removes use of Safe.pm altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include that it is now possible to use Perl’s strict pragma in a natural way in plperl, and that Perl’s $a and $b variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

• Prevent PL/Tcl from executing untrustworthy code from pltcl_modules (Tom)

PL/Tcl’s feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless pltcl_modules is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted “normal” Tcl interpreter unless we are really going to execute a pltclu function. (CVE-2010-1170)

• Fix possible crash if a cache reset message is received during rebuild of a relcache entry (Heikki)

This error was introduced in 8.2.16 while fixing a related failure.

• Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

Previously, if an unprivileged user ran ALTER USER ... RESET ALL for himself, or ALTER DATABASE ... RESET ALL for a database he owns, this would remove all special parameter
settings for the user or database, even ones that are only supposed to be changeable by a superuser. Now, the ALTER will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a CONTEXT addition would be made to log entries (Tom)

In some cases the context-printing function would fail because the current transaction had already been rolled back when it came time to print a log message.

- Update PL/Perl's ppport.h for modern Perl versions (Andrew)
- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)
- Prevent infinite recursion in psql when expanding a variable that refers to itself (Tom)
- Fix psql's \copy to not add spaces around a dot within \copy (select ...) (Tom)
  
  Addition of spaces around the decimal point in a numeric literal would result in a syntax error.
- Ensure that contrib/pgstattuple functions respond to cancel interrupts promptly (Tatsuhito Kasahara)
- Make server startup deal properly with the case that shmget() returns EINVAL for an existing shared memory segment (Tom)

  This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entirely-misleading startup failure complaining that the shared memory request size was too large.
- Avoid possible crashes in syslogger process on Windows (Heikki)
- Deal more robustly with incomplete time zone information in the Windows registry (Magnus)
- Update the set of known Windows time zone names (Magnus)
- Update time zone data files to tzdata release 2010j for DST law changes in Argentina, Australian Antarctic, Bangladesh, Mexico, Morocco, Pakistan, Palestine, Russia, Syria, Tunisia; also historical corrections for Taiwan.

  Also, add PKST (Pakistan Summer Time) to the default set of timezone abbreviations.

**E.213. Release 8.2.16**

**Release date:** 2010-03-15

This release contains a variety of fixes from 8.2.15. For information about new features in the 8.2 major release, see Section E.229.

**E.213.1. Migration to Version 8.2.16**

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

**E.213.2. Changes**

- Add new configuration parameter ssl_renegotiation_limit to control how often we do session key renegotiation for an SSL connection (Magnus)

  This can be set to zero to disable renegotiation completely, which may be required if a broken SSL library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that cause renegotiation attempts to fail.
• Fix possible deadlock during backend startup (Tom)

• Fix possible crashes due to not handling errors during relcache reload cleanly (Tom)

• Fix possible crashes when trying to recover from a failure in subtransaction start (Tom)

• Fix server memory leak associated with use of savepoints and a client encoding different from server's encoding (Tom)

• Fix incorrect WAL data emitted during end-of-recovery cleanup of a GIST index page split (Yoichi Hirai)

  This would result in index corruption, or even more likely an error during WAL replay, if we were unlucky enough to crash during end-of-recovery cleanup after having completed an incomplete GIST insertion.

• Make substring() for bit types treat any negative length as meaning “all the rest of the string” (Tom)

  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

• Fix integer-to-bit-string conversions to handle the first fractional byte correctly when the output bit width is wider than the given integer by something other than a multiple of 8 bits (Tom)

• Fix some cases of pathologically slow regular expression matching (Tom)

• Fix the STOP WAL LOCATION entry in backup history files to report the next WAL segment's name when the end location is exactly at a segment boundary (Itagaki Takahiro)

• Fix some more cases of temporary-file leakage (Heikki)

  This corrects a problem introduced in the previous minor release. One case that failed is when a plpgsql function returning set is called within another function's exception handler.

• Improve constraint exclusion processing of boolean-variable cases, in particular make it possible to exclude a partition that has a “bool_column = false” constraint (Tom)

• When reading pg_hba.conf and related files, do not treat @something as a file inclusion request if the @ appears inside quote marks; also, never treat @ by itself as a file inclusion request (Tom)

  This prevents erratic behavior if a role or database name starts with @. If you need to include a file whose path name contains spaces, you can still do so, but you must write @"/path to/file" rather than putting the quotes around the whole construct.

• Prevent infinite loop on some platforms if a directory is named as an inclusion target in pg_hba.conf and related files (Tom)

• Fix possible infinite loop if SSL_read or SSL_write fails without setting errno (Tom)

  This is reportedly possible with some Windows versions of OpenSSL.

• Fix psql's numericlocale option to not format strings it shouldn't in latex and troff output formats (Heikki)

• Make psql return the correct exit status (3) when ON_ERROR_STOP and --single-transaction are both specified and an error occurs during the implied COMMIT (Bruce)

• Fix plpgsql failure in one case where a composite column is set to NULL (Tom)

• Fix possible failure when calling PL/Perl functions from PL/PerlU or vice versa (Tim Bunce)
• Add volatile markings in PL/Python to avoid possible compiler-specific misbehavior (Zdenek Kotala)

• Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)

The only known symptom of this oversight is that the Tcl clock command misbehaves if using Tcl 8.5 or later.

• Prevent crash in contrib/dblink when too many key columns are specified to a dblink_build_sql_* function (Rushabh Lathia, Joe Conway)

• Fix assorted crashes in contrib/xml2 caused by sloppy memory management (Tom)

• Make building of contrib/xml2 more robust on Windows (Andrew)

• Fix race condition in Windows signal handling (Radu Ilie)

One known symptom of this bug is that rows in pg_listener could be dropped under heavy load.

• Update time zone data files to tzdata release 2010e for DST law changes in Bangladesh, Chile, Fiji, Mexico, Paraguay, Samoa.

E.214. Release 8.2.15

Release date: 2009-12-14

This release contains a variety of fixes from 8.2.14. For information about new features in the 8.2 major release, see Section E.229.

E.214.1. Migration to Version 8.2.15

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.14, see Section E.215.

E.214.2. Changes

• Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)

This change prevents allegedly-immutable index functions from possibly subverting a superuser's session (CVE-2009-4136).

• Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)

This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).

• Fix possible crash during backend-startup-time cache initialization (Tom)

• Prevent signals from interrupting VACUUM at unsafe times (Alvaro)

This fix prevents a PANIC if a VACUUM FULL is canceled after it's already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

• Fix possible crash due to integer overflow in hash table size calculation (Tom)

This could occur with extremely large planner estimates for the size of a hashjoin's result.

• Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)
• Ensure that shared tuple-level locks held by prepared transactions are not ignored (Heikki)

• Fix premature drop of temporary files used for a cursor that is accessed within a subtransaction (Heikki)

• Fix incorrect logic for GiST index page splits, when the split depends on a non-first column of the index (Paul Ramsey)

• Don’t error out if recycling or removing an old WAL file fails at the end of checkpoint (Heikki)

   It’s better to treat the problem as non-fatal and allow the checkpoint to complete. Future checkpoints will retry the removal. Such problems are not expected in normal operation, but have been seen to be caused by misdesigned Windows anti-virus and backup software.

• Ensure WAL files aren’t repeatedly archived on Windows (Heikki)

   This is another symptom that could happen if some other process interfered with deletion of a no-longer-needed file.

• Fix PAM password processing to be more robust (Tom)

   The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.

• Fix processing of ownership dependencies during CREATE OR REPLACE FUNCTION (Tom)

• Fix bug with calling plperl from plperlu or vice versa (Tom)

   An error exit from the inner function could result in crashes due to failure to re-select the correct Perl interpreter for the outer function.

• Fix session-lifespan memory leak when a PL/Perl function is redefined (Tom)

• Ensure that Perl arrays are properly converted to PostgreSQL arrays when returned by a set-returning PL/Perl function (Andrew Dunstan, Abhijit Menon-Sen)

   This worked correctly already for non-set-returning functions.

• Fix rare crash in exception processing in PL/Python (Peter)

• Ensure psql's flex module is compiled with the correct system header definitions (Tom)

   This fixes build failures on platforms where --enable-largefile causes incompatible changes in the generated code.

• Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

• Update the timezone abbreviation files to match current reality (Joachim Wieland)

   This includes adding IDT and SGT to the default timezone abbreviation set.

• Update time zone data files to tzdata release 2009s for DST law changes in Antarctica, Argentina, Bangladesh, Fiji, Novokuznetsk, Pakistan, Palestine, Samoa, Syria; also historical corrections for Hong Kong.

E.215. Release 8.2.14

Release date: 2009-09-09
This release contains a variety of fixes from 8.2.13. For information about new features in the 8.2 major release, see Section E.229.

E.215.1. Migration to Version 8.2.14

A dump/restore is not required for those running 8.2.X. However, if you have any hash indexes on interval columns, you must REINDEX them after updating to 8.2.14. Also, if you are upgrading from a version earlier than 8.2.11, see Section E.218.

E.215.2. Changes

- Force WAL segment switch during `pg_start_backup()` (Heikki)
  This avoids corner cases that could render a base backup unusable.

- Disallow `RESET ROLE` and `RESET SESSION AUTHORIZATION` inside security-definer functions (Tom, Heikki)
  This covers a case that was missed in the previous patch that disallowed `SET ROLE` and `SET SESSION AUTHORIZATION` inside security-definer functions. (See CVE-2007-6600)

- Make `LOAD` of an already-loaded loadable module into a no-op (Tom)
  Formerly, `LOAD` would attempt to unload and re-load the module, but this is unsafe and not all that useful.

- Disallow empty passwords during LDAP authentication (Magnus)

- Fix handling of sub-SELECTs appearing in the arguments of an outer-level aggregate function (Tom)

- Fix bugs associated with fetching a whole-row value from the output of a Sort or Materialize plan node (Tom)

- Revert planner change that disabled partial-index and constraint exclusion optimizations when there were more than 100 clauses in an AND or OR list (Tom)

- Fix hash calculation for data type `interval` (Tom)
  This corrects wrong results for hash joins on interval values. It also changes the contents of hash indexes on interval columns. If you have any such indexes, you must REINDEX them after updating.

- Treat `to_char(..., 'TH')` as an uppercase ordinal suffix with 'HH'/HH12' (Heikki)
  It was previously handled as 'th' (lowercase).

- Fix overflow for `INTERVAL 'x ms'` when x is more than 2 million and integer datetimes are in use (Alex Hunsaker)

- Fix calculation of distance between a point and a line segment (Tom)
  This led to incorrect results from a number of geometric operators.

- Fix `money` data type to work in locales where currency amounts have no fractional digits, e.g. Japan (Itagaki Takahiro)

- Properly round datetime input like `00:12:57.9999999999999999999999999999` (Tom)

- Fix poor choice of page split point in GiST R-tree operator classes (Teodor)
• Avoid performance degradation in bulk inserts into GIN indexes when the input values are (nearly) in sorted order (Tom)

• Correctly enforce NOT NULL domain constraints in some contexts in PL/pgSQL (Tom)

• Fix portability issues in plperl initialization (Andrew Dunstan)

• Fix pg_ctl to not go into an infinite loop if postgresql.conf is empty (Jeff Davis)

• Make contrib/hstore throw an error when a key or value is too long to fit in its data structure, rather than silently truncating it (Andrew Gierth)

• Fix contrib/xml2's xslt_process() to properly handle the maximum number of parameters (twenty) (Tom)

• Improve robustness of libpq's code to recover from errors during COPY FROM STDIN (Tom)

• Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)

• Update time zone data files to tzdata release 2009l for DST law changes in Bangladesh, Egypt, Jordan, Pakistan, Argentina/San_Luis, Cuba, Jordan (historical correction only), Mauritius, Morocco, Palestine, Syria, Tunisia.

E.216. Release 8.2.13

Release date: 2009-03-16

This release contains a variety of fixes from 8.2.12. For information about new features in the 8.2 major release, see Section E.229.

E.216.1. Migration to Version 8.2.13

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.11, see Section E.218.

E.216.2. Changes

• Prevent error recursion crashes when encoding conversion fails (Tom)

  This change extends fixes made in the last two minor releases for related failure scenarios. The previous fixes were narrowly tailored for the original problem reports, but we have now recognized that any error thrown by an encoding conversion function could potentially lead to infinite recursion while trying to report the error. The solution therefore is to disable translation and encoding conversion and report the plain-ASCII form of any error message, if we find we have gotten into a recursive error reporting situation. (CVE-2009-0922)

• Disallow CREATE CONVERSION with the wrong encodings for the specified conversion function (Heikki)

  This prevents one possible scenario for encoding conversion failure. The previous change is a backstop to guard against other kinds of failures in the same area.

• Fix core dump when to_char() is given format codes that are inappropriate for the type of the data argument (Tom)

• Fix possible failure in contrib/tsearch2 when C locale is used with a multi-byte encoding (Teodor)

  Crashes were possible on platforms where wchar_t is narrower than int; Windows in particular.
• Fix extreme inefficiency in contrib/tsearch2 parser's handling of an email-like string containing multiple @ characters (Heikki)

• Fix decompilation of CASE WHEN with an implicit coercion (Tom)

This mistake could lead to Assert failures in an Assert-enabled build, or an “unexpected CASE WHEN clause” error message in other cases, when trying to examine or dump a view.

• Fix possible misassignment of the owner of a TOAST table's rowtype (Tom)

If CLUSTER or a rewriting variant of ALTER TABLE were executed by someone other than the table owner, the pg_type entry for the table's TOAST table would end up marked as owned by that someone. This caused no immediate problems, since the permissions on the TOAST rowtype aren't examined by any ordinary database operation. However, it could lead to unexpected failures if one later tried to drop the role that issued the command (in 8.1 or 8.2), or “owner of data type appears to be invalid” warnings from pg_dump after having done so (in 8.3).

• Fix PL/pgSQL to not treat INTO after INSERT as an INTO-variables clause anywhere in the string, not only at the start; in particular, don’t fail for INSERT INTO within CREATE RULE (Tom)

• Clean up PL/pgSQL error status variables fully at block exit (Ashesh Vashi and Dave Page)

This is not a problem for PL/pgSQL itself, but the omission could cause the PL/pgSQL Debugger to crash while examining the state of a function.

• Retry failed calls to CallNamedPipe() on Windows (Steve Marshall, Magnus)

It appears that this function can sometimes fail transiently; we previously treated any failure as a hard error, which could confuse LISTEN/NOTIFY as well as other operations.

• Add MUST (Mauritius Island Summer Time) to the default list of known timezone abbreviations (Xavier Bugaud)

### E.217. Release 8.2.12

**Release date:** 2009-02-02

This release contains a variety of fixes from 8.2.11. For information about new features in the 8.2 major release, see Section E.229.

### E.217.1. Migration to Version 8.2.12

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.11, see Section E.218.

### E.217.2. Changes

• Improve handling of URLs in headline() function (Teodor)

• Improve handling of overlength headlines in headline() function (Teodor)

• Prevent possible Assert failure or misconversion if an encoding conversion is created with the wrong conversion function for the specified pair of encodings (Tom, Heikki)

• Fix possible Assert failure if a statement executed in PL/pgSQL is rewritten into another kind of statement, for example if an INSERT is rewritten into an UPDATE (Heikki)

• Ensure that a snapshot is available to datatype input functions (Tom)
This primarily affects domains that are declared with CHECK constraints involving user-defined stable or immutable functions. Such functions typically fail if no snapshot has been set.

- Make it safer for SPI-using functions to be used within datatype I/O; in particular, to be used in domain check constraints (Tom)

- Avoid unnecessary locking of small tables in VACUUM (Heikki)

- Fix a problem that made UPDATE RETURNING tableoid return zero instead of the correct OID (Tom)

- Fix planner misestimation of selectivity when transitive equality is applied to an outer-join clause (Tom)

  This could result in bad plans for queries like ... from a left join b on a.a1 = b.b1 where a.a1 = 42 ...

- Improve optimizer's handling of long IN lists (Tom)

  This change avoids wasting large amounts of time on such lists when constraint exclusion is enabled.

- Ensure that the contents of a holdable cursor don't depend on the contents of TOAST tables (Tom)

  Previously, large field values in a cursor result might be represented as TOAST pointers, which would fail if the referenced table got dropped before the cursor is read, or if the large value is deleted and then vacuumed away. This cannot happen with an ordinary cursor, but it could with a cursor that is held past its creating transaction.

- Fix memory leak when a set-returning function is terminated without reading its whole result (Tom)

- Fix contrib/dblink's dblink_get_result(text,bool) function (Joe)

- Fix possible garbage output from contrib/sslinfo functions (Tom)

- Fix configure script to properly report failure when unable to obtain linkage information for PL/Perl (Andrew)

  - Make all documentation reference pgsql-bugs and/or pgsql-hackers as appropriate, instead of the now-decommissioned pgsql-ports and pgsql-patches mailing lists (Tom)

  - Update time zone data files to tzdata release 2009a (for Kathmandu and historical DST corrections in Switzerland, Cuba)

E.218. Release 8.2.11

Release date: 2008-11-03

This release contains a variety of fixes from 8.2.10. For information about new features in the 8.2 major release, see Section E.229.

E.218.1. Migration to Version 8.2.11

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.7, see Section E.222. Also, if you were running a previous 8.2.X release, it is recommended to REINDEX all GiST indexes after the upgrade.

E.218.2. Changes

- Fix GiST index corruption due to marking the wrong index entry “dead” after a deletion (Teodor)
This would result in index searches failing to find rows they should have found. Corrupted indexes can be fixed with `REINDEX`.

- **Fix backend crash when the client encoding cannot represent a localized error message (Tom)**

  We have addressed similar issues before, but it would still fail if the “character has no equivalent” message itself couldn't be converted. The fix is to disable localization and send the plain ASCII error message when we detect such a situation.

- **Fix possible crash when deeply nested functions are invoked from a trigger (Tom)**

- **Improve optimization of `expression IN(expression-list)` queries (Tom, per an idea from Robert Haas)**

  Cases in which there are query variables on the right-hand side had been handled less efficiently in 8.2.x and 8.3.x than in prior versions. The fix restores 8.1 behavior for such cases.

- **Fix mis-expansion of rule queries when a sub-SELECT appears in a function call in FROM, a multi-row VALUES list, or a RETURNING list (Tom)**

  The usual symptom of this problem is an “unrecognized node type” error.

- **Fix memory leak during rescan of a hashed aggregation plan (Neil)**

- **Ensure an error is reported when a newly-defined PL/pgSQL trigger function is invoked as a normal function (Tom)**

- **Prevent possible collision of relfilenode numbers when moving a table to another tablespace with `ALTER SET TABLESPACE` (Heikki)**

  The command tried to re-use the existing filename, instead of picking one that is known unused in the destination directory.

- **Fix incorrect tsearch2 headline generation when single query item matches first word of text (Sushant Sinha)**

- **Fix improper display of fractional seconds in interval values when using a non-ISO datestyle in an `--enable-integer-datetimes` build (Ron Mayer)**

- **Ensure `SPI_getvalue` and `SPI_getbinval` behave correctly when the passed tuple and tuple descriptor have different numbers of columns (Tom)**

  This situation is normal when a table has had columns added or removed, but these two functions didn't handle it properly. The only likely consequence is an incorrect error indication.

- **Fix ecpg's parsing of `CREATE ROLE` (Michael)**

- **Fix recent breakage of `pg_ctl restart` (Tom)**

- **Ensure `pg_control` is opened in binary mode (Itagaki Takahiro)**

  `pg_controldata` and `pg_resetxlog` did this incorrectly, and so could fail on Windows.

- **Update time zone data files to tzdata release 2008i (for DST law changes in Argentina, Brazil, Mauritius, Syria)**

E.219. Release 8.2.10

**Release date:** 2008-09-22
This release contains a variety of fixes from 8.2.9. For information about new features in the 8.2 major release, see Section E.229.

E.219.1. Migration to Version 8.2.10

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.7, see Section E.222.

E.219.2. Changes

- Fix bug in btree WAL recovery code (Heikki)
  Recovery failed if the WAL ended partway through a page split operation.
- Fix potential miscalculation of `datfrozenxid` (Alvaro)
  This error may explain some recent reports of failure to remove old pg_clog data.
- Widen local lock counters from 32 to 64 bits (Tom)
  This responds to reports that the counters could overflow in sufficiently long transactions, leading to unexpected “lock is already held” errors.
- Fix possible duplicate output of tuples during a GiST index scan (Teodor)
- Fix missed permissions checks when a view contains a simple `UNION ALL` construct (Heikki)
  Permissions for the referenced tables were checked properly, but not permissions for the view itself.
- Add checks in executor startup to ensure that the tuples produced by an `INSERT` or `UPDATE` will match the target table’s current rowtype (Tom)
  `ALTER COLUMN TYPE`, followed by re-use of a previously cached plan, could produce this type of situation. The check protects against data corruption and/or crashes that could ensue.
- Fix possible repeated drops during `DROP OWNED` (Tom)
  This would typically result in strange errors such as “cache lookup failed for relation NNN”.
- Fix `AT TIME ZONE` to first try to interpret its timezone argument as a timezone abbreviation, and only try it as a full timezone name if that fails, rather than the other way around as formerly (Tom)
  The timestamp input functions have always resolved ambiguous zone names in this order. Making `AT TIME ZONE` do so as well improves consistency, and fixes a compatibility bug introduced in 8.1: in ambiguous cases we now behave the same as 8.0 and before did, since in the older versions `AT TIME ZONE` accepted only abbreviations.
- Fix datetime input functions to correctly detect integer overflow when running on a 64-bit platform (Tom)
- Prevent integer overflows during units conversion when displaying a configuration parameter that has units (Tom)
- Improve performance of writing very long log messages to syslog (Tom)
- Allow spaces in the suffix part of an LDAP URL in `pg_hba.conf` (Tom)
- Fix bug in backwards scanning of a cursor on a `SELECT DISTINCT ON` query (Tom)
- Fix planner bug with nested sub-select expressions (Tom)
If the outer sub-select has no direct dependency on the parent query, but the inner one does, the outer value might not get recalculated for new parent query rows.

- Fix planner to estimate that \texttt{GROUP BY} expressions yielding boolean results always result in two groups, regardless of the expressions' contents (Tom)

  This is very substantially more accurate than the regular \texttt{GROUP BY} estimate for certain boolean tests like \texttt{col IS NULL}.

- Fix PL/pgSQL to not fail when a \texttt{FOR} loop's target variable is a record containing composite-type fields (Tom)

- Fix PL/Tcl to behave correctly with Tcl 8.5, and to be more careful about the encoding of data sent to or from Tcl (Tom)

- On Windows, work around a Microsoft bug by preventing libpq from trying to send more than 64kB per system call (Magnus)

- Improve \texttt{pg_dump} and \texttt{pg_restore}'s error reporting after failure to send a SQL command (Tom)

- Fix \texttt{pg_ctl} to properly preserve postmaster command-line arguments across a restart (Bruce)

- Update time zone data files to \texttt{tzdata} release 2008f (for DST law changes in Argentina, Bahamas, Brazil, Mauritius, Morocco, Pakistan, Palestine, and Paraguay)

\section*{E.220. Release 8.2.9}

\textbf{Release date:} 2008-06-12

This release contains one serious and one minor bug fix over 8.2.8. For information about new features in the 8.2 major release, see Section E.229.

\section*{E.220.1. Migration to Version 8.2.9}

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.7, see Section E.222.

\section*{E.220.2. Changes}

- Make \texttt{pg\_get\_ruledef()} parenthesize negative constants (Tom)

  Before this fix, a negative constant in a view or rule might be dumped as, say, \texttt{-42::integer}, which is subtly incorrect: it should be \texttt{(-42)::integer} due to operator precedence rules. Usually this would make little difference, but it could interact with another recent patch to cause PostgreSQL to reject what had been a valid \texttt{SELECT DISTINCT} view query. Since this could result in \texttt{pg\_dump} output failing to reload, it is being treated as a high-priority fix. The only released versions in which dump output is actually incorrect are 8.3.1 and 8.2.7.

- Make \texttt{ALTER AGGREGATE ... OWNER TO} update \texttt{pg\_shdepend} (Tom)

  This oversight could lead to problems if the aggregate was later involved in a \texttt{DROP OWNED} or \texttt{REASSIGN OWNED} operation.

\section*{E.221. Release 8.2.8}

\textbf{Release date:} never released
This release contains a variety of fixes from 8.2.7. For information about new features in the 8.2 major release, see Section E.229.

E.221.1. Migration to Version 8.2.8

A dump/restore is not required for those running 8.2.X. However, if you are upgrading from a version earlier than 8.2.7, see Section E.222.

E.221.2. Changes

- Fix ERRORDATA_STACK_SIZE exceeded crash that occurred on Windows when using UTF-8 database encoding and a different client encoding (Tom)

- Fix ALTER TABLE ADD COLUMN ... PRIMARY KEY so that the new column is correctly checked to see if it's been initialized to all non-nulls (Brendan Jurd)

  Previous versions neglected to check this requirement at all.

- Fix possible CREATE TABLE failure when inheriting the “same” constraint from multiple parent relations that inherited that constraint from a common ancestor (Tom)

- Fix pg_get_ruledef() to show the alias, if any, attached to the target table of an UPDATE or DELETE (Tom)

- Fix GIN bug that could result in a too many LWLocks taken failure (Teodor)

- Avoid possible crash when decompressing corrupted data (Zdenek Kotala)

- Repair two places where SIGTERM exit of a backend could leave corrupted state in shared memory (Tom)

  Neither case is very important if SIGTERM is used to shut down the whole database cluster together, but there was a problem if someone tried to SIGTERM individual backends.

- Fix conversions between ISO-8859-5 and other encodings to handle Cyrillic “Yo” characters (e and E with two dots) (Sergey Burladyan)

- Fix several datatype input functions, notably array_in(), that were allowing unused bytes in their results to contain uninitialized, unpredictable values (Tom)

  This could lead to failures in which two apparently identical literal values were not seen as equal, resulting in the parser complaining about unmatched ORDER BY and DISTINCT expressions.

- Fix a corner case in regular-expression substring matching (substring(string from pattern)) (Tom)

  The problem occurs when there is a match to the pattern overall but the user has specified a parenthesized subexpression and that subexpression hasn't got a match. An example is substring('foo' from 'foo(bar)?'). This should return NULL, since (bar) isn't matched, but it was mistakenly returning the whole-pattern match instead (ie, foo).

- Update time zone data files to tzdata release 2008c (for DST law changes in Morocco, Iraq, Choibalsan, Pakistan, Syria, Cuba, and Argentina/San_Luis)

- Fix incorrect result from ecpg's PGTYPEStimestamp_sub() function (Michael)

- Fix broken GiST comparison function for contrib/tsearch2's tsquery type (Teodor)

- Fix possible crashes in contrib/cube functions (Tom)
• Fix core dump in contrib/xml2’s xpath_table() function when the input query returns a NULL value (Tom)

• Fix contrib/xml2’s makefile to not override CFLAGS (Tom)

• Fix DatumGetBool macro to not fail with gcc 4.3 (Tom)

This problem affects “old style” (V0) C functions that return boolean. The fix is already in 8.3, but the need to back-patch it was not realized at the time.

E.222. Release 8.2.7

Release date: 2008-03-17

This release contains a variety of fixes from 8.2.6. For information about new features in the 8.2 major release, see Section E.229.

E.222.1. Migration to Version 8.2.7

A dump/restore is not required for those running 8.2.X. However, you might need to REINDEX indexes on textual columns after updating, if you are affected by the Windows locale issue described below.

E.222.2. Changes

• Fix character string comparison for Windows locales that consider different character combinations as equal (Tom)

  This fix applies only on Windows and only when using UTF-8 database encoding. The same fix was made for all other cases over two years ago, but Windows with UTF-8 uses a separate code path that was not updated. If you are using a locale that considers some non-identical strings as equal, you may need to REINDEX to fix existing indexes on textual columns.

• Repair potential deadlock between concurrent VACUUM FULL operations on different system catalogs (Tom)

• Fix longstanding LISTEN/NOTIFY race condition (Tom)

  In rare cases a session that had just executed a LISTEN might not get a notification, even though one would be expected because the concurrent transaction executing NOTIFY was observed to commit later.

  A side effect of the fix is that a transaction that has executed a not-yet-committed LISTEN command will not see any row in pg_listener for the LISTEN, should it choose to look; formerly it would have. This behavior was never documented one way or the other, but it is possible that some applications depend on the old behavior.

• Disallow LISTEN and UNLISTEN within a prepared transaction (Tom)

  This was formerly allowed but trying to do it had various unpleasant consequences, notably that the originating backend could not exit as long as an UNLISTEN remained uncommitted.

• Disallow dropping a temporary table within a prepared transaction (Heikki)

  This was correctly disallowed by 8.1, but the check was inadvertently broken in 8.2.

• Fix rare crash when an error occurs during a query using a hash index (Heikki)

• Fix memory leaks in certain usages of set-returning functions (Neil)
• Fix input of datetime values for February 29 in years BC (Tom)
  The former coding was mistaken about which years were leap years.
• Fix “unrecognized node type” error in some variants of ALTER OWNER (Tom)
• Ensure pg_stat_activity.waiting flag is cleared when a lock wait is aborted (Tom)
• Fix handling of process permissions on Windows Vista (Dave, Magnus)
  In particular, this fix allows starting the server as the Administrator user.
• Update time zone data files to tzdata release 2008a (in particular, recent Chile changes); adjust timezone abbreviation VET (Venezuela) to mean UTC-4:30, not UTC-4:00 (Tom)
• Fix pg_ctl to correctly extract the postmaster's port number from command-line options (Itagaki Takahiro, Tom)
  Previously, pg_ctl start -w could try to contact the postmaster on the wrong port, leading to bogus reports of startup failure.
• Use -fwrapv to defend against possible misoptimization in recent gcc versions (Tom)
  This is known to be necessary when building PostgreSQL with gcc 4.3 or later.
• Correctly enforce statement_timeout values longer than INT_MAX microseconds (about 35 minutes) (Tom)
  This bug affects only builds with --enable-integer-datetimes.
• Fix “unexpected PARAM_SUBLINK ID” planner error when constant-folding simplifies a sub-select (Tom)
• Fix logical errors in constraint-exclusion handling of IS NULL and NOT expressions (Tom)
  The planner would sometimes exclude partitions that should not have been excluded because of the possibility of NULL results.
• Fix another cause of “failed to build any N-way joins” planner errors (Tom)
  This could happen in cases where a clauseless join needed to be forced before a join clause could be exploited.
• Fix incorrect constant propagation in outer-join planning (Tom)
  The planner could sometimes incorrectly conclude that a variable could be constrained to be equal to a constant, leading to wrong query results.
• Fix display of constant expressions in ORDER BY and GROUP BY (Tom)
  An explicitly casted constant would be shown incorrectly. This could for example lead to corruption of a view definition during dump and reload.
• Fix libpq to handle NOTICE messages correctly during COPY OUT (Tom)
  This failure has only been observed to occur when a user-defined datatype's output routine issues a NOTICE, but there is no guarantee it couldn't happen due to other causes.

E.223. Release 8.2.6

Release date: 2008-01-07
This release contains a variety of fixes from 8.2.5, including fixes for significant security issues. For information about new features in the 8.2 major release, see Section E.229.

### E.223.1. Migration to Version 8.2.6

A dump/restore is not required for those running 8.2.X.

### E.223.2. Changes

- Prevent functions in indexes from executing with the privileges of the user running `VACUUM`, `ANALYZE`, etc (Tom)

  Functions used in index expressions and partial-index predicates are evaluated whenever a new table entry is made. It has long been understood that this poses a risk of trojan-horse code execution if one modifies a table owned by an untrustworthy user. (Note that triggers, defaults, check constraints, etc. pose the same type of risk.) But functions in indexes pose extra danger because they will be executed by routine maintenance operations such as `VACUUM FULL`, which are commonly performed automatically under a superuser account. For example, a nefarious user can execute code with superuser privileges by setting up a trojan-horse index definition and waiting for the next routine vacuum. The fix arranges for standard maintenance operations (including `VACUUM`, `ANALYZE`, `REINDEX`, and `CLUSTER`) to execute as the table owner rather than the calling user, using the same privilege-switching mechanism already used for `SECURITY DEFINER` functions. To prevent bypassing this security measure, execution of `SET SESSION AUTHORIZATION` and `SET ROLE` is now forbidden within a `SECURITY DEFINER` context. (CVE-2007-6600)

- Repair assorted bugs in the regular-expression package (Tom, Will Drewry)

  Suitably crafted regular-expression patterns could cause crashes, infinite or near-infinite looping, and/or massive memory consumption, all of which pose denial-of-service hazards for applications that accept regex search patterns from untrustworthy sources. (CVE-2007-4769, CVE-2007-4772, CVE-2007-6067)

- Require non-superusers who use `/contrib/dblink` to use only password authentication, as a security measure (Joe)

  The fix that appeared for this in 8.2.5 was incomplete, as it plugged the hole for only some `dblink` functions. (CVE-2007-6601, CVE-2007-3278)

- Fix bugs in WAL replay for GIN indexes (Teodor)

- Fix GIN index build to work properly when `maintenance_work_mem` is 4GB or more (Tom)

- Update time zone data files to tzdata release 2007k (in particular, recent Argentina changes) (Tom)

- Improve planner’s handling of LIKE/regex estimation in non-C locales (Tom)

- Fix planning-speed problem for deep outer-join nests, as well as possible poor choice of join order (Tom)

- Fix planner failure in some cases of `WHERE false AND var IN (SELECT ...)` (Tom)

- Make `CREATE TABLE ... SERIAL` and `ALTER SEQUENCE ... OWNED BY` not change the `currval()` state of the sequence (Tom)

- Preserve the tablespace and storage parameters of indexes that are rebuilt by `ALTER TABLE ... ALTER COLUMN TYPE` (Tom)

- Make archive recovery always start a new WAL timeline, rather than only when a recovery stop time was used (Simon)
This avoids a corner-case risk of trying to overwrite an existing archived copy of the last WAL segment, and seems simpler and cleaner than the original definition.

- Make VACUUM not use all of maintenance_work_mem when the table is too small for it to be useful (Alvaro)
- Fix potential crash in translate() when using a multibyte database encoding (Tom)
- Make corr() return the correct result for negative correlation values (Neil)
- Fix overflow in extract(epoch from interval) for intervals exceeding 68 years (Tom)
- Fix PL/Perl to not fail when a UTF-8 regular expression is used in a trusted function (Andrew)
- Fix PL/Perl to cope when platform's Perl defines type bool as int rather than char (Tom)
  While this could theoretically happen anywhere, no standard build of Perl did things this way ... until macOS 10.5.
- Fix PL/Python to work correctly with Python 2.5 on 64-bit machines (Marko Kreen)
- Fix PL/Python to not crash on long exception messages (Alvaro)
- Fix pg_dump to correctly handle inheritance child tables that have default expressions different from their parent's (Tom)
- Fix libpq crash when PGPASSFILE refers to a file that is not a plain file (Martin Pitt)
- ecpg parser fixes (Michael)
- Make contrib/pgcrypto defend against OpenSSL libraries that fail on keys longer than 128 bits; which is the case at least on some Solaris versions (Marko Kreen)
- Make contrib/tablefunc's crosstab() handle NULL rowid as a category in its own right, rather than crashing (Joe)
- Fix tsvector and tsquery output routines to escape backslashes correctly (Teodor, Bruce)
- Fix crash of to_tsvector() on huge input strings (Teodor)
- Require a specific version of Autoconf to be used when re-generating the configure script (Peter)
  This affects developers and packagers only. The change was made to prevent accidental use of untested combinations of Autoconf and PostgreSQL versions. You can remove the version check if you really want to use a different Autoconf version, but it's your responsibility whether the result works or not.
- Update gettimeofday configuration check so that PostgreSQL can be built on newer versions of MinGW (Magnus)

### E.224. Release 8.2.5

**Release date:** 2007-09-17

This release contains a variety of fixes from 8.2.4. For information about new features in the 8.2 major release, see Section E.229.

### E.224.1. Migration to Version 8.2.5

A dump/restore is not required for those running 8.2.X.
E.224.2. Changes

- Prevent index corruption when a transaction inserts rows and then aborts close to the end of a concurrent VACUUM on the same table (Tom)

- Fix ALTER DOMAIN ADD CONSTRAINT for cases involving domains over domains (Tom)

- Make CREATE DOMAIN ... DEFAULT NULL work properly (Tom)

- Fix some planner problems with outer joins, notably poor size estimation for t1 LEFT JOIN t2 WHERE t2.col IS NULL (Tom)

- Allow the interval data type to accept input consisting only of milliseconds or microseconds (Neil)

- Allow timezone name to appear before the year in timestamp input (Tom)

- Fixes for GIN indexes used by /contrib/tsearch2 (Teodor)

- Speed up rtree index insertion (Teodor)

- Fix excessive logging of SSL error messages (Tom)

- Fix logging so that log messages are never interleaved when using the syslogger process (Andrew)

- Fix crash when log_min_error_statement logging runs out of memory (Tom)

- Fix incorrect handling of some foreign-key corner cases (Tom)

- Fix stddev_pop(numeric) and var_pop(numeric) (Tom)

- Prevent REINDEX and CLUSTER from failing due to attempting to process temporary tables of other sessions (Alvaro)

- Update the time zone database rules, particularly New Zealand's upcoming changes (Tom)

- Windows socket and semaphore improvements (Magnus)

- Make pg_ctl -w work properly in Windows service mode (Dave Page)

- Fix memory allocation bug when using MIT Kerberos on Windows (Magnus)

- Suppress timezone name (%Z) in log timestamps on Windows because of possible encoding mismatches (Tom)

- Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

- Restrict /contrib/pgstattuple functions to superusers, for security reasons (Tom)

- Do not let /contrib/intarray try to make its GIN opclass the default (this caused problems at dump/restore) (Tom)

E.225. Release 8.2.4

Release date: 2007-04-23

This release contains a variety of fixes from 8.2.3, including a security fix. For information about new features in the 8.2 major release, see Section E.229.
E.225.1. Migration to Version 8.2.4

A dump/restore is not required for those running 8.2.X.

E.225.2. Changes

- Support explicit placement of the temporary-table schema within search_path, and disable searching it for functions and operators (Tom)

  This is needed to allow a security-definer function to set a truly secure value of search_path. Without it, an unprivileged SQL user can use temporary objects to execute code with the privileges of the security-definer function (CVE-2007-2138). See CREATE FUNCTION for more information.

- Fix shared_preload_libraries for Windows by forcing reload in each backend (Korry Douglas)

- Fix to_char() so it properly upper/lower cases localized day or month names (Pavel Stehule)

- /contrib/tsearch2 crash fixes (Teodor)

- Require COMMIT PREPARED to be executed in the same database as the transaction was prepared in (Heikki)

- Allow pg_dump to do binary backups larger than two gigabytes on Windows (Magnus)

- New traditional (Taiwan) Chinese FAQ (Zhou Daojing)

- Prevent the statistics collector from writing to disk too frequently (Tom)

- Fix potential-data-corruption bug in how VACUUM FULL handles UPDATE chains (Tom, Pavan Deolasee)

- Fix bug in domains that use array types (Tom)

- Fix pg_dump so it can dump a serial column's sequence using -t when not also dumping the owning table (Tom)

- Planner fixes, including improving outer join and bitmap scan selection logic (Tom)

- Fix possible wrong answers or crash when a PL/pgSQL function tries to RETURN from within an EXCEPTION block (Tom)

- Fix PANIC during enlargement of a hash index (Tom)

- Fix POSIX-style timezone specs to follow new USA DST rules (Tom)

E.226. Release 8.2.3

Release date: 2007-02-07

This release contains two fixes from 8.2.2. For information about new features in the 8.2 major release, see Section E.229.

E.226.1. Migration to Version 8.2.3

A dump/restore is not required for those running 8.2.X.

E.226.2. Changes

- Remove overly-restrictive check for type length in constraints and functional indexes (Tom)
• Fix optimization so MIN/MAX in subqueries can again use indexes (Tom)

E.227. Release 8.2.2

Release date: 2007-02-05

This release contains a variety of fixes from 8.2.1, including a security fix. For information about new features in the 8.2 major release, see Section E.229.

E.227.1. Migration to Version 8.2.2

A dump/restore is not required for those running 8.2.X.

E.227.2. Changes

• Remove security vulnerabilities that allowed connected users to read backend memory (Tom)

  The vulnerabilities involve suppressing the normal check that a SQL function returns the data type it's declared to, and changing the data type of a table column (CVE-2007-0555, CVE-2007-0556). These errors can easily be exploited to cause a backend crash, and in principle might be used to read database content that the user should not be able to access.

• Fix not-so-rare-anymore bug wherein btree index page splits could fail due to choosing an infeasible split point (Heikki Linnakangas)

• Fix Borland C compile scripts (L Bayuk)

• Properly handle `to_char('CC')` for years ending in 00 (Tom)

  Year 2000 is in the twentieth century, not the twenty-first.

• /contrib/tsearch2 localization improvements (Tatsuo, Teodor)

• Fix incorrect permission check in `information_schema.key_column_usage` view (Tom)

  The symptom is "relation with OID nnnnn does not exist" errors. To get this fix without using `initdb`, use `CREATE OR REPLACE VIEW` to install the corrected definition found in `share/information_schema.sql`. Note you will need to do this in each database.

• Improve VACUUM performance for databases with many tables (Tom)

• Fix for rare Assert() crash triggered by `UNION` (Tom)

• Fix potentially incorrect results from index searches using `ROW` inequality conditions (Tom)

• Tighten security of multi-byte character processing for UTF8 sequences over three bytes long (Tom)

• Fix bogus "permission denied" failures occurring on Windows due to attempts to fsync already-deleted files (Magnus, Tom)

• Fix bug that could cause the statistics collector to hang on Windows (Magnus)

  This would in turn lead to autovacuum not working.

• Fix possible crashes when an already-in-use PL/pgSQL function is updated (Tom)

• Improve PL/pgSQL handling of domain types (Sergiy Vyshnevetskiy, Tom)

• Fix possible errors in processing PL/pgSQL exception blocks (Tom)
E.228. Release 8.2.1

Release date: 2007-01-08

This release contains a variety of fixes from 8.2. For information about new features in the 8.2 major release, see Section E.229.

E.228.1. Migration to Version 8.2.1

A dump/restore is not required for those running 8.2.

E.228.2. Changes

- Fix crash with `SELECT ... LIMIT ALL` (also `LIMIT NULL`) (Tom)
- Several `/contrib/tsearch2` fixes (Teodor)
- On Windows, make log messages coming from the operating system use ASCII encoding (Hiroshi Saito)
  
  This fixes a conversion problem when there is a mismatch between the encoding of the operating system and database server.
- Fix Windows linking of `pg_dump` using `win32.mak` (Hiroshi Saito)
- Fix planner mistakes for outer join queries (Tom)
- Fix several problems in queries involving sub-SELECTs (Tom)
- Fix potential crash in SPI during subtransaction abort (Tom)
  
  This affects all PL functions since they all use SPI.
- Improve build speed of PDF documentation (Peter)
- Re-add JST (Japan) timezone abbreviation (Tom)
- Improve optimization decisions related to index scans (Tom)
- Have `psql` print multi-byte combining characters as before, rather than output as `\u` (Tom)
- Improve index usage of regular expressions that use parentheses (Tom)
  
  This improves `psql` command performance also.
- Make `pg_dumpall` assume that databases have public `CONNECT` privilege, when dumping from a pre-8.2 server (Tom)
  
  This preserves the previous behavior that anyone can connect to a database if allowed by `pg_hba.conf`.

E.229. Release 8.2

Release date: 2006-12-05

E.229.1. Overview

This release adds many functionality and performance improvements that were requested by users, including:
• Query language enhancements including INSERT/UPDATE/DELETE RETURNING, multirow VALUES lists, and optional target-table alias in UPDATE/DELETE

• Index creation without blocking concurrent INSERT/UPDATE/DELETE operations

• Many query optimization improvements, including support for reordering outer joins

• Improved sorting performance with lower memory usage

• More efficient locking with better concurrency

• More efficient vacuuming

• Easier administration of warm standby servers

• New FILLFACTOR support for tables and indexes

• Monitoring, logging, and performance tuning additions

• More control over creating and dropping objects

• Table inheritance relationships can be defined for and removed from pre-existing tables

• COPY TO can copy the output of an arbitrary SELECT statement

• Array improvements, including nulls in arrays

• Aggregate-function improvements, including multiple-input aggregates and SQL:2003 statistical functions

• Many contrib/ improvements

## E.229.2. Migration to Version 8.2

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

• Set escape_string_warning to on by default (Bruce)

  This issues a warning if backslash escapes are used in non-escape (non-''') strings.

• Change the row constructor syntax (ROW(...)) so that list elements foo.* will be expanded to a list of their member fields, rather than creating a nested row type field as formerly (Tom)

  The new behavior is substantially more useful since it allows, for example, triggers to check for data changes with IF row(new.*) IS DISTINCT FROM row(old.*). The old behavior is still available by omitting .*.

• Make row comparisons follow SQL standard semantics and allow them to be used in index scans (Tom)

  Previously, row = and <> comparisons followed the standard but < <= > >= did not. A row comparison can now be used as an index constraint for a multicolumn index matching the row value.

• Make row IS [NOT] NULL tests follow SQL standard semantics (Tom)

  The former behavior conformed to the standard for simple cases with IS NULL, but IS NOT NULL would return true if any row field was non-null, whereas the standard says it should return true only when all fields are non-null.

• Make SET CONSTRAINT affect only one constraint (Kris Jurka)
In previous releases, SET CONSTRAINT modified all constraints with a matching name. In this release, the schema search path is used to modify only the first matching constraint. A schema specification is also supported. This more nearly conforms to the SQL standard.

- **Remove RULE permission for tables, for security reasons (Tom)**

  As of this release, only a table's owner can create or modify rules for the table. For backwards compatibility, GRANT/REVOKE RULE is still accepted, but it does nothing.

- **Array comparison improvements (Tom)**

  Now array dimensions are also compared.

- **Change array concatenation to match documented behavior (Tom)**

  This changes the previous behavior where concatenation would modify the array lower bound.

- **Make command-line options of postmaster and postgres identical (Peter)**

  This allows the postmaster to pass arguments to each backend without using -o. Note that some options are now only available as long-form options, because there were conflicting single-letter options.

- **Deprecate use of postmaster symbolic link (Peter)**

  postmaster and postgres commands now act identically, with the behavior determined by command-line options. The postmaster symbolic link is kept for compatibility, but is not really needed.

- **Change log_duration to output even if the query is not output (Tom)**

  In prior releases, log_duration only printed if the query appeared earlier in the log.

- **Make to_char(time) and to_char(interval) treat HH and HH12 as 12-hour intervals**

  Most applications should use HH24 unless they want a 12-hour display.

- **Zero unmasked bits in conversion from INET to CIDR (Tom)**

  This ensures that the converted value is actually valid for CIDR.

- **Remove australian_timezones configuration variable (Joachim Wieland)**

  This variable has been superseded by a more general facility for configuring timezone abbreviations.

- **Improve cost estimation for nested-loop index scans (Tom)**

  This might eliminate the need to set unrealistically small values of random_page_cost. If you have been using a very small random_page_cost, please recheck your test cases.

- **Change behavior of pg_dump --n and --t options. (Greg Sabino Mullane)**

  See the pg_dump manual page for details.

- **Change libpq PQdsplen() to return a useful value (Martijn van Oosterhout)**

- **Declare libpq PQgetssl() as returning void *, rather than SSL * (Martijn van Oosterhout)**

  This allows applications to use the function without including the OpenSSL headers.

- **C-language loadable modules must now include a PG_MODULE_MAGIC macro call for version compatibility checking (Martijn van Oosterhout)**
• For security’s sake, modules used by a PL/PerlU function are no longer available to PL/Perl functions (Andrew)

Note
This also implies that data can no longer be shared between a PL/Perl function and a PL/PerlU function. Some Perl installations have not been compiled with the correct flags to allow multiple interpreters to exist within a single process. In this situation PL/Perl and PL/PerlU cannot both be used in a single backend. The solution is to get a Perl installation which supports multiple interpreters.

• In contrib/xml2/, rename xml_valid() to xml_is_well_formed() (Tom)
xml_valid() will remain for backward compatibility, but its behavior will change to do schema checking in a future release.

• Remove contrib/ora2pg/, now at http://www.samse.fr/GPL/ora2pg

• Remove contrib modules that have been migrated to PgFoundry: addepend, dbase, dbmirror, fulltextindex, mac, userlock

• Remove abandoned contrib modules: mSQL-interface, tips

• Remove QNX and BEOS ports (Bruce)
These ports no longer had active maintainers.

E.229.3. Changes
Below you will find a detailed account of the changes between PostgreSQL 8.2 and the previous major release.

E.229.3.1. Performance Improvements

• Allow the planner to reorder outer joins in some circumstances (Tom)
In previous releases, outer joins would always be evaluated in the order written in the query. This change allows the query optimizer to consider reordering outer joins, in cases where it can determine that the join order can be changed without altering the meaning of the query. This can make a considerable performance difference for queries involving multiple outer joins or mixed inner and outer joins.

• Improve efficiency of IN (list-of-expressions) clauses (Tom)

• Improve sorting speed and reduce memory usage (Simon, Tom)

• Improve subtransaction performance (Alvaro, Itagaki Takahiro, Tom)

• Add FILLFACTOR to table and index creation (ITAGAKI Takahiro)
This leaves extra free space in each table or index page, allowing improved performance as the database grows. This is particularly valuable to maintain clustering.

• Increase default values for shared_buffers and max_fsm_pages (Andrew)

• Improve locking performance by breaking the lock manager tables into sections (Tom)
This allows locking to be more fine-grained, reducing contention.
• Reduce locking requirements of sequential scans (Qingqing Zhou)

• Reduce locking required for database creation and destruction (Tom)

• Improve the optimizer's selectivity estimates for LIKE, ILIKE, and regular expression operations (Tom)

• Improve planning of joins to inherited tables and UNION ALL views (Tom)

• Allow constraint exclusion to be applied to inherited UPDATE and DELETE queries (Tom)
  
  SELECT already honored constraint exclusion.

• Improve planning of constant WHERE clauses, such as a condition that depends only on variables inherited from an outer query level (Tom)

• Protocol-level unnamed prepared statements are re-planned for each set of BIND values (Tom)
  This improves performance because the exact parameter values can be used in the plan.

• Speed up vacuuming of B-Tree indexes (Heikki Linnakangas, Tom)

• Avoid extra scan of tables without indexes during VACUUM (Greg Stark)

• Improve multicolumn GiST indexing (Oleg, Teodor)

• Remove dead index entries before B-Tree page split (Junji Teramoto)

E.229.3.2. Server Changes

• Allow a forced switch to a new transaction log file (Simon, Tom)
  
  This is valuable for keeping warm standby slave servers in sync with the master. Transaction log file switching now also happens automatically during pg_stop_backup(). This ensures that all transaction log files needed for recovery can be archived immediately.

• Add WAL informational functions (Simon)
  
  Add functions for interrogating the current transaction log insertion point and determining WAL filenames from the hex WAL locations displayed by pg_stop_backup() and related functions.

• Improve recovery from a crash during WAL replay (Simon)
  
  The server now does periodic checkpoints during WAL recovery, so if there is a crash, future WAL recovery is shortened. This also eliminates the need for warm standby servers to replay the entire log since the base backup if they crash.

• Improve reliability of long-term WAL replay (Heikki, Simon, Tom)
  
  Formerly, trying to roll forward through more than 2 billion transactions would not work due to XID wraparound. This meant warm standby servers had to be reloaded from fresh base backups periodically.

• Add archive_timeout to force transaction log file switches at a given interval (Simon)
  
  This enforces a maximum replication delay for warm standby servers.

• Add native LDAP authentication (Magnus Hagander)
  
  This is particularly useful for platforms that do not support PAM, such as Windows.

• Add GRANT CONNECT ON DATABASE (Gevik Babakhani)
This gives SQL-level control over database access. It works as an additional filter on top of the existing `pg_hba.conf` controls.

- Add support for SSL Certificate Revocation List (CRL) files (Libor Hohoš)
  The server and libpq both recognize CRL files now.
- GiST indexes are now clusterable (Teodor)
- Remove routine autovacuum server log entries (Bruce)
  `pg_stat_activity` now shows autovacuum activity.
- Track maximum XID age within individual tables, instead of whole databases (Alvaro)
  This reduces the overhead involved in preventing transaction ID wraparound, by avoiding unnecessary VACUUMs.
- Add last vacuum and analyze timestamp columns to the stats collector (Larry Rosenman)
  These values now appear in the `pg_stat_*_tables` system views.
- Improve performance of statistics monitoring, especially `stats_command_string` (Tom, Bruce)
  This release enables `stats_command_string` by default, now that its overhead is minimal. This means `pg_stat_activity` will now show all active queries by default.
- Add a `waiting` column to `pg_stat_activity` (Tom)
  This allows `pg_stat_activity` to show all the information included in the `ps` display.
- Add configuration parameter `update_process_title` to control whether the `ps` display is updated for every command (Bruce)
  On platforms where it is expensive to update the `ps` display, it might be worthwhile to turn this off and rely solely on `pg_stat_activity` for status information.
- Allow units to be specified in configuration settings (Peter)
  For example, you can now set `shared_buffers` to 32MB rather than mentally converting sizes.
- Add support for include directives in `postgresql.conf` (Joachim Wieland)
- Improve logging of protocol-level prepare/bind/execute messages (Bruce, Tom)
  Such logging now shows statement names, bind parameter values, and the text of the query being executed. Also, the query text is properly included in logged error messages when enabled by `log_min_error_statement`.
- Prevent `max_stack_depth` from being set to unsafe values
  On platforms where we can determine the actual kernel stack depth limit (which is most), make sure that the initial default value of `max_stack_depth` is safe, and reject attempts to set it to unsafely large values.
- Enable highlighting of error location in query in more cases (Tom)
  The server is now able to report a specific error location for some semantic errors (such as unrecognized column name), rather than just for basic syntax errors as before.
- Fix “failed to re-find parent key” errors in `VACUUM` (Tom)
Release Notes

- Clean out pg_internal.init cache files during server restart (Simon)
  
  This avoids a hazard that the cache files might contain stale data after PITR recovery.
- Fix race condition for truncation of a large relation across a gigabyte boundary by VACUUM (Tom)
- Fix bug causing needless deadlock errors on row-level locks (Tom)
- Fix bugs affecting multi-gigabyte hash indexes (Tom)
- Each backend process is now its own process group leader (Tom)

  This allows query cancel to abort subprocesses invoked from a backend or archive/recovery process.

E.229.3.3. Query Changes

- Add INSERT/UPDATE/DELETE RETURNING (Jonah Harris, Tom)

  This allows these commands to return values, such as the computed serial key for a new row. In the UPDATE case, values from the updated version of the row are returned.
- Add support for multiple-row VALUES clauses, per SQL standard (Joe, Tom)

  This allows INSERT to insert multiple rows of constants, or queries to generate result sets using constants. For example, INSERT ... VALUES (...), (...), ..... and SELECT * FROM (VALUES (...), (...), ...) AS alias(f1, ...).
- Allow UPDATE and DELETE to use an alias for the target table (Atsushi Ogawa)

  The SQL standard does not permit an alias in these commands, but many database systems allow one anyway for notational convenience.
- Allow UPDATE to set multiple columns with a list of values (Susanne Ebrecht)

  This is basically a short-hand for assigning the columns and values in pairs. The syntax is UPDATE tab SET (column, ...) = (val, ...).
- Make row comparisons work per standard (Tom)

  The forms <, <=, >, >= now compare rows lexicographically, that is, compare the first elements, if equal compare the second elements, and so on. Formerly they expanded to an AND condition across all the elements, which was neither standard nor very useful.
- Add CASCADE option to TRUNCATE (Joachim Wieland)

  This causes TRUNCATE to automatically include all tables that reference the specified table(s) via foreign keys. While convenient, this is a dangerous tool — use with caution!
- Support FOR UPDATE and FOR SHARE in the same SELECT command (Tom)
- Add IS NOT DISTINCT FROM (Pavel Stehule)

  This operator is similar to equality (=), but evaluates to true when both left and right operands are NULL, and to false when just one is, rather than yielding NULL in these cases.
- Improve the length output used by UNION/INTERSECT/EXCEPT (Tom)

  When all corresponding columns are of the same defined length, that length is used for the result, rather than a generic length.
- Allow ILIKE to work for multi-byte encodings (Tom)
Internally, ILIKE now calls lower() and then uses LIKE. Locale-specific regular expression patterns still do not work in these encodings.

- Enable standard_conforming_strings to be turned on (Kevin Grittner)
  
  This allows backslash escaping in strings to be disabled, making PostgreSQL more standards-compliant. The default is off for backwards compatibility, but future releases will default this to on.

- Do not flatten subqueries that contain volatile functions in their target lists (Jaime Casanova)
  
  This prevents surprising behavior due to multiple evaluation of a volatile function (such as random() or nextval()). It might cause performance degradation in the presence of functions that are unnecessarily marked as volatile.

- Add system views pg_prepared_statements and pg_cursors to show prepared statements and open cursors (Joachim Wieland, Neil)

  These are very useful in pooled connection setups.

- Support portal parameters in EXPLAIN and EXECUTE (Tom)

  This allows, for example, JDBC parameters to work in these commands.

- If SQL-level PREPARE parameters are unspecified, infer their types from the content of the query (Neil)

  Protocol-level PREPARE already did this.

- Allow LIMIT and OFFSET to exceed two billion (Dhanaraj M)

### E.229.3.4. Object Manipulation Changes

- Add TABLESPACE clause to CREATE TABLE AS (Neil)

  This allows a tablespace to be specified for the new table.

- Add ON COMMIT clause to CREATE TABLE AS (Neil)

  This allows temporary tables to be truncated or dropped on transaction commit. The default behavior is for the table to remain until the session ends.

- Add INCLUDING CONSTRAINTS to CREATE TABLE LIKE (Greg Stark)

  This allows easy copying of CHECK constraints to a new table.

- Allow the creation of placeholder (shell) types (Martijn van Oosterhout)

  A shell type declaration creates a type name, without specifying any of the details of the type. Making a shell type is useful because it allows cleaner declaration of the type’s input/output functions, which must exist before the type can be defined “for real”. The syntax is CREATE TYPE type-name.

- Aggregate functions now support multiple input parameters (Sergey Koposov, Tom)

- Add new aggregate creation syntax (Tom)

  The new syntax is CREATE AGGREGATE agname (input_type) (parameter_list).

  This more naturally supports the new multi-parameter aggregate functionality. The previous syntax is still supported.

- Add ALTER ROLE PASSWORD NULL to remove a previously set role password (Peter)
• Add DROP object IF EXISTS for many object types (Andrew)
  This allows DROP operations on non-existent objects without generating an error.

• Add DROP OWNED to drop all objects owned by a role (Alvaro)

• Add REASSIGN OWNED to reassign ownership of all objects owned by a role (Alvaro)
  This, and DROP OWNED above, facilitate dropping roles.

• Add GRANT ON SEQUENCE syntax (Bruce)
  This was added for setting sequence-specific permissions. GRANT ON TABLE for sequences is
  still supported for backward compatibility.

• Add USAGE permission for sequences that allows only currval() and nextval(), not setval() (Bruce)
  USAGE permission allows more fine-grained control over sequence access. Granting USAGE allows
  users to increment a sequence, but prevents them from setting the sequence to an arbitrary value
  using setval().

• Add ALTER TABLE [ NO ] INHERIT (Greg Stark)
  This allows inheritance to be adjusted dynamically, rather than just at table creation and destruction.
  This is very valuable when using inheritance to implement table partitioning.

• Allow comments on global objects to be stored globally (Kris Jurka)
  Previously, comments attached to databases were stored in individual databases, making them ineffectual, and there was no provision at all for comments on roles or tablespaces. This change adds a new shared catalog pg_shdescription and stores comments on databases, roles, and tablespaces therein.

E.229.3.5. Utility Command Changes

• Add option to allow indexes to be created without blocking concurrent writes to the table (Greg Stark, Tom)
  The new syntax is CREATE INDEX CONCURRENTLY. The default behavior is still to block table
  modification while an index is being created.

• Provide advisory locking functionality (Abhijit Menon-Sen, Tom)
  This is a new locking API designed to replace what used to be in /contrib/userlock. The userlock
  code is now on pgfoundry.

• Allow COPY to dump a SELECT query (Zoltan Boszormenyi, Karel Zak)
  This allows COPY to dump arbitrary SQL queries. The syntax is COPY (SELECT ...) TO.

• Make the COPY command return a command tag that includes the number of rows copied (Volkan YAZICI)

• Allow VACUUM to expire rows without being affected by other concurrent VACUUM operations
  (Hannu Krossing, Alvaro, Tom)

• Make initdb detect the operating system locale and set the default DateStyle accordingly (Peter)
  This makes it more likely that the installed postgresql.conf DateStyle value will be as desired.

• Reduce number of progress messages displayed by initdb (Tom)
E.229.3.6. Date/Time Changes

- Allow full timezone names in timestamp input values (Joachim Wieland)
  For example, '2006-05-24 21:11 America/New_York':timestamptz.
- Support configurable timezone abbreviations (Joachim Wieland)
  A desired set of timezone abbreviations can be chosen via the configuration parameter timezone_abbreviations.
- Add pg_timezone_abbrs and pg_timezone_names views to show supported timezones (Magnus Hagander)
- Add clock_timestamp(), statement_timestamp(), and transaction_timestamp() (Bruce)
  clock_timestamp() is the current wall-clock time, statement_timestamp() is the time the current statement arrived at the server, and transaction_timestamp() is an alias for now().
- Allow to_char() to print localized month and day names (Euler Taveira de Oliveira)
- Allow to_char(time) and to_char(interval) to output AM/PM specifications (Bruce)
  Intervals and times are treated as 24-hour periods, e.g. 25 hours is considered AM.
- Add new function justify_interval() to adjust interval units (Mark Dilger)
- Allow timezone offsets up to 14:59 away from GMT
  Kiribati uses GMT+14, so we’d better accept that.
- Interval computation improvements (Michael GlAESemann, Bruce)

E.229.3.7. Other Data Type and Function Changes

- Allow arrays to contain NULL elements (Tom)
- Allow assignment to array elements not contiguous with the existing entries (Tom)
  The intervening array positions will be filled with nulls. This is per SQL standard.
- New built-in operators for array-subset comparisons (@>, <@, &&) (Teodor, Tom)
  These operators can be indexed for many data types using GiST or GIN indexes.
- Add convenient arithmetic operations on INET/CIDR values (Stephen R. van den Berg)
  The new operators are & (and), | (or), ~ (not), inet + int8, inet - int8, and inet - inet.
- Add new aggregate functions from SQL:2003 (Neil)
  The new functions are var_pop(), var_samp(), stddev_pop(), and stddev_samp().
  var_samp() and stddev_samp() are merely renamings of the existing aggregates variance() and stddev(). The latter names remain available for backward compatibility.
- Add SQL:2003 statistical aggregates (Sergey Koposov)
  New functions: regr_intercept(), regr_slope(), regr_r2(), corr(), covar_samp(), covar_pop(), regr_avgx(), regr_avgy(), regr_sxy(), regr_sxx(), regr_syy(), regr_count().
• Allow domains to be based on other domains (Tom)
• Properly enforce domain CHECK constraints everywhere (Neil, Tom)

For example, the result of a user-defined function that is declared to return a domain type is now checked against the domain's constraints. This closes a significant hole in the domain implementation.

• Fix problems with dumping renamed SERIAL columns (Tom)

The fix is to dump a SERIAL column by explicitly specifying its DEFAULT and sequence elements, and reconstructing the SERIAL column on reload using a new ALTER SEQUENCE OWNED BY command. This also allows dropping a SERIAL column specification.

• Add a server-side sleep function pg_sleep() (Joachim Wieland)

• Add all comparison operators for the tid (tuple id) data type (Mark Kirkwood, Greg Stark, Tom)

E.229.3.8. PL/pgSQL Server-Side Language Changes

• Add TG_table_name and TG_table_schema to trigger parameters (Andrew)

TG_relname is now deprecated. Comparable changes have been made in the trigger parameters for the other PLs as well.

• Allow FOR statements to return values to scalars as well as records and row types (Pavel Stehule)
• Add a BY clause to the FOR loop, to control the iteration increment (Jaime Casanova)
• Add STRICT to SELECT INTO (Matt Miller)

STRICT mode throws an exception if more or less than one row is returned by the SELECT, for Oracle PL/SQL compatibility.

E.229.3.9. PL/Perl Server-Side Language Changes

• Add table_name and table_schema to trigger parameters (Adam Sjøgren)

• Add prepared queries (Dmitry Karasik)
• Make $_TD trigger data a global variable (Andrew)

Previously, it was lexical, which caused unexpected sharing violations.

• Run PL/Perl and PL/PerlU in separate interpreters, for security reasons (Andrew)

In consequence, they can no longer share data nor loaded modules. Also, if Perl has not been compiled with the requisite flags to allow multiple interpreters, only one of these languages can be used in any given backend process.

E.229.3.10. PL/Python Server-Side Language Changes

• Named parameters are passed as ordinary variables, as well as in the args[] array (Sven Suursoho)
• Add table_name and table_schema to trigger parameters (Andrew)
• Allow returning of composite types and result sets (Sven Suursoho)
• Return result-set as list, iterator, or generator (Sven Suursoho)
• Allow functions to return void (Neil)
• Python 2.5 is now supported (Tom)

E.229.3.11. psql Changes

• Add new command \password for changing role password with client-side password encryption (Peter)
• Allow \c to connect to a new host and port number (David, Volkan YAZICI)
• Add tablespace display to \l+ (Philip Yarra)
• Improve \df slash command to include the argument names and modes (OUT or INOUT) of the function (David Fetter)
• Support binary COPY (Andreas Pflug)
• Add option to run the entire session in a single transaction (Simon)
  Use option -1 or --single-transaction.
• Support for automatically retrieving SELECT results in batches using a cursor (Chris Mair)
  This is enabled using \set FETCH_COUNT n. This feature allows large result sets to be retrieved in psql without attempting to buffer the entire result set in memory.
• Make multi-line values align in the proper column (Martijn van Oosterhout)
  Field values containing newlines are now displayed in a more readable fashion.
• Save multi-line statements as a single entry, rather than one line at a time (Sergey E. Koposov)
  This makes up-arrow recall of queries easier. (This is not available on Windows, because that platform uses the native command-line editing present in the operating system.)
• Make the line counter 64-bit so it can handle files with more than two billion lines (David Fetter)
• Report both the returned data and the command status tag for INSERT/UPDATE/DELETE RETURNING (Tom)

E.229.3.12. pg_dump Changes

• Allow complex selection of objects to be included or excluded by pg_dump (Greg Sabino Mullane)
  pg_dump now supports multiple -n (schema) and -t (table) options, and adds -N and -T options to exclude objects. Also, the arguments of these switches can now be wild-card expressions rather than single object names, for example -t 'foo*', and a schema can be part of a -t or -T switch, for example -t schema1.table1.
• Add pg_restore --no-data-for-failed-tables option to suppress loading data if table creation failed (i.e., the table already exists) (Martin Pitt)
• Add pg_restore option to run the entire session in a single transaction (Simon)
  Use option -1 or --single-transaction.

E.229.3.13. libpq Changes

• Add PQencryptPassword() to encrypt passwords (Tom)
  This allows passwords to be sent pre-encrypted for commands like ALTER ROLE ... PASSWORD.
• Add function PQisthreadsafe() (Bruce)
  This allows applications to query the thread-safety status of the library.

• Add PQdescribePrepared(), PQdescribePortal(), and related functions to return information about previously prepared statements and open cursors (Volkan YAZICI)

• Allow LDAP lookups from pg_service.conf (Laurenz Albe)

• Allow a hostname in ~/.pgpass to match the default socket directory (Bruce)
  A blank hostname continues to match any Unix-socket connection, but this addition allows entries that are specific to one of several postmasters on the machine.

E.229.3.14. ecpg Changes

• Allow SHOW to put its result into a variable (Joachim Wieland)

• Add COPY TO STDOUT (Joachim Wieland)

• Add regression tests (Joachim Wieland, Michael)

• Major source code cleanups (Joachim Wieland, Michael)

E.229.3.15. Windows Port

• Allow MSVC to compile the PostgreSQL server (Magnus, Hiroshi Saito)

• Add MSVC support for utility commands and pg_dump (Hiroshi Saito)

• Add support for Windows code pages 1253, 1254, 1255, and 1257 (Kris Jurka)

• Drop privileges on startup, so that the server can be started from an administrative account (Magnus)

• Stability fixes (Qingqing Zhou, Magnus)

• Add native semaphore implementation (Qingqing Zhou)
  The previous code mimicked SysV semaphores.

E.229.3.16. Source Code Changes

• Add GIN (Generalized Inverted iNdex) index access method (Teodor, Oleg)

• Remove R-tree indexing (Tom)
  Rtree has been re-implemented using GiST. Among other differences, this means that rtree indexes now have support for crash recovery via write-ahead logging (WAL).

• Reduce libraries needlessly linked into the backend (Martijn van Oosterhout, Tom)

• Add a configure flag to allow libedit to be preferred over GNU readline (Bruce)
  Use configure --with-libedit-preferred.

• Allow installation into directories containing spaces (Peter)

• Improve ability to relocate installation directories (Tom)

• Add support for Solaris x86_64 using the Solaris compiler (Pierre Girard, Theo Schlossnagle, Bruce)

• Add DTrace support (Robert Lor)

3079
• Add **PG_VERSION_NUM** for use by third-party applications wanting to test the backend version in C using > and < comparisons (Bruce)

• Add **XLOG_BLCKSZ** as independent from **BLCKSZ** (Mark Wong)

• Add **LWLOCK_STATS** define to report locking activity (Tom)

• Emit warnings for unknown configure options (Martijn van Oosterhout)

• Add server support for “plugin” libraries that can be used for add-on tasks such as debugging and performance measurement (Korry Douglas)

This consists of two features: a table of “rendezvous variables” that allows separately-loaded shared libraries to communicate, and a new configuration parameter **local_preload_libraries** that allows libraries to be loaded into specific sessions without explicit cooperation from the client application. This allows external add-ons to implement features such as a PL/pgSQL debugger.

• Rename existing configuration parameter **preload_libraries** to **shared_preload_libraries** (Tom)

This was done for clarity in comparison to **local_preload_libraries**.

• Add new configuration parameter **server_version_num** (Greg Sabino Mullane)

This is like **server_version**, but is an integer, e.g. 80200. This allows applications to make version checks more easily.

• Add a configuration parameter **seq_page_cost** (Tom)

• Re-implement the regression test script as a C program (Magnus, Tom)

• Allow loadable modules to allocate shared memory and lightweight locks (Marc Munro)

• Add automatic initialization and finalization of dynamically loaded libraries (Ralf Engelschall, Tom)

New functions **_PG_init()** and **_PG_fini()** are called if the library defines such symbols. Hence we no longer need to specify an initialization function in **shared_preload_libraries**; we can assume that the library used the **_PG_init()** convention instead.

• Add **PG_MODULE_MAGIC** header block to all shared object files (Martijn van Oosterhout)

The magic block prevents version mismatches between loadable object files and servers.

• Add shared library support for AIX (Laurenz Albe)

• New XML documentation section (Bruce)

**E.229.3.17. Contrib Changes**

• Major tsearch2 improvements (Oleg, Teodor)

  • multibyte encoding support, including UTF8
  
  • query rewriting support
  
  • improved ranking functions
  
  • thesaurus dictionary support
  
  • Ispell dictionaries now recognize MySpell format, used by OpenOffice
    
  • GIN support
• Add adminpack module containing Pgadmin administration functions (Dave)
  These functions provide additional file system access routines not present in the default PostgreSQL server.

• Add sslinfo module (Victor Wagner)
  Reports information about the current connection’s SSL certificate.

• Add pgrowlocks module (Tatsuo)
  This shows row locking information for a specified table.

• Add hstore module (Oleg, Teodor)

• Add isn module, replacing isbn_isbn (Jeremy Kronuz)
  This new implementation supports EAN13, UPC, ISBN (books), ISMN (music), and ISSN (serials).

• Add index information functions to pgstattriple (ITAGAKI Takahiro, Satoshi Nagayasu)

• Add pg_freespacemap module to display free space map information (Mark Kirkwood)

• pgcrypto now has all planned functionality (Marko Kreen)
  • Include iMath library in pgcrypto to have the public-key encryption functions always available.
  • Add SHA224 algorithm that was missing in OpenBSD code.
  • Activate builtin code for SHA224/256/384/512 hashes on older OpenSSL to have those algorithms always available.
  • New function gen_random_bytes() that returns cryptographically strong randomness. Useful for generating encryption keys.
  • Remove digest_exists(), hmac_exists() and cipher_exists() functions.

• Improvements to cube module (Joshua Reich)
  New functions are cube(float[]), cube(float[], float[]), and cube_subset(cube, int4[]).

• Add async query capability to dblink (Kai Londenberg, Joe Conway)

• New operators for array-subset comparisons (@>, <@, &@) (Tom)
  Various contrib packages already had these operators for their datatypes, but the naming wasn't consistent. We have now added consistently named array-subset comparison operators to the core code and all the contrib packages that have such functionality. (The old names remain available, but are deprecated.)

• Add uninstall scripts for all contrib packages that have install scripts (David, Josh Drake)

E.230. Release 8.1.23

Release date: 2010-12-16

This release contains a variety of fixes from 8.1.22. For information about new features in the 8.1 major release, see Section E.253.

This is expected to be the last PostgreSQL release in the 8.1.X series. Users are encouraged to update to a newer release branch soon.
E.230.1. Migration to Version 8.1.23

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.18, see Section E.235.

E.230.2. Changes

• Force the default `wal_sync_method` to be `fdatasync` on Linux (Tom Lane, Marti Raudsepp)

The default on Linux has actually been `fdatasync` for many years, but recent kernel changes caused PostgreSQL to choose `open_dataasync` instead. This choice did not result in any performance improvement, and caused outright failures on certain filesystems, notably `ext4` with the `data=journal` mount option.

• Fix recovery from base backup when the starting checkpoint WAL record is not in the same WAL segment as its redo point (Jeff Davis)

• Add support for detecting register-stack overrun on IA64 (Tom Lane)

The IA64 architecture has two hardware stacks. Full prevention of stack-overrun failures requires checking both.

• Add a check for stack overflow in `copyObject()` (Tom Lane)

Certain code paths could crash due to stack overflow given a sufficiently complex query.

• Fix detection of page splits in temporary GiST indexes (Heikki Linnakangas)

It is possible to have a “concurrent” page split in a temporary index, if for example there is an open cursor scanning the index when an insertion is done. GiST failed to detect this case and hence could deliver wrong results when execution of the cursor continued.

• Avoid memory leakage while `ANALYZE`’ing complex index expressions (Tom Lane)

• Ensure an index that uses a whole-row Var still depends on its table (Tom Lane)

An index declared like `create index i on t (foo(t.*))` would not automatically get dropped when its table was dropped.

• Do not “inline” a SQL function with multiple `OUT` parameters (Tom Lane)

This avoids a possible crash due to loss of information about the expected result rowtype.

• Fix constant-folding of `COALESCE()` expressions (Tom Lane)

The planner would sometimes attempt to evaluate sub-expressions that in fact could never be reached, possibly leading to unexpected errors.

• Add print functionality for `InhRelation` nodes (Tom Lane)

This avoids a failure when `debug_print_parse` is enabled and certain types of query are executed.

• Fix incorrect calculation of distance from a point to a horizontal line segment (Tom Lane)

This bug affected several different geometric distance-measurement operators.

• Fix PL/pgSQL’s handling of “simple” expressions to not fail in recursion or error-recovery cases (Tom Lane)

• Fix bug in `contrib/cube`’s GiST picksplit algorithm (Alexander Korotkov)
This could result in considerable inefficiency, though not actually incorrect answers, in a GiST index on a cube column. If you have such an index, consider REINDEXing it after installing this update.

- Don’t emit “identifier will be truncated” notices in contrib/dblink except when creating new connections (Itagaki Takahiro)
- Fix potential coredump on missing public key in contrib/pgcrypto (Marti Raudsepp)
- Fix memory leak in contrib/xml2’s XPath query functions (Tom Lane)
- Update time zone data files to t z data release 2010o for DST law changes in Fiji and Samoa; also historical corrections for Hong Kong.

E.231. Release 8.1.22

Release date: 2010-10-04

This release contains a variety of fixes from 8.1.21. For information about new features in the 8.1 major release, see Section E.253.

The PostgreSQL community will stop releasing updates for the 8.1.X release series in November 2010. Users are encouraged to update to a newer release branch soon.

E.231.1. Migration to Version 8.1.22

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.18, see Section E.235.

E.231.2. Changes

- Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

  This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function’s owner.

  The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

  It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

  Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

- Prevent possible crashes in pg_get_expr() by disallowing it from being called with an argument that is not one of the system catalog columns it’s intended to be used with (Heikki Linnakangas, Tom Lane)

- Fix “cannot handle unplanned sub-select” error (Tom Lane)

  This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.
• Prevent show_session_authorization() from crashing within autovacuum processes (Tom Lane)

• Defend against functions returning setof record where not all the returned rows are actually of the same rowtype (Tom Lane)

• Fix possible failure when hashing a pass-by-reference function result (Tao Ma, Tom Lane)

• Take care to fsync the contents of lockfiles (both postmaster.pid and the socket lockfile) while writing them (Tom Lane)

  This omission could result in corrupted lockfile contents if the machine crashes shortly after postmaster start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.

• Avoid recursion while assigning XIDs to heavily-nested subtransactions (Andres Freund, Robert Haas)

  The original coding could result in a crash if there was limited stack space.

• Fix log_line_prefix's %i escape, which could produce junk early in backend startup (Tom Lane)

• Fix possible data corruption in ALTER TABLE ... SET TABLESPACE when archiving is enabled (Jeff Davis)

• Allow CREATE DATABASE and ALTER DATABASE ... SET TABLESPACE to be interrupted by query-cancel (Guillaume Lelarge)

• In PL/Python, defend against null pointer results from PyObject_AsVoidPtr and PyObject_FromVoidPtr (Peter Eisentraut)

• Improve contrib/dblink's handling of tables containing dropped columns (Tom Lane)

• Fix connection leak after “duplicate connection name” errors in contrib/dblink (Itagaki Takahiro)

• Fix contrib/dblink to handle connection names longer than 62 bytes correctly (Itagaki Takahiro)

• Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)

• Update time zone data files to tzdata release 2010l for DST law changes in Egypt and Palestine; also historical corrections for Finland.

  This change also adds new names for two Micronesian timezones: Pacific/Chuuk is now preferred over Pacific/Truk (and the preferred abbreviation is CHUT not TRUT) and Pacific/Pohnpei is preferred over Pacific/Ponape.

E.232. Release 8.1.21

Release date: 2010-05-17

This release contains a variety of fixes from 8.1.20. For information about new features in the 8.1 major release, see Section E.253.

E.232.1. Migration to Version 8.1.21

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.18, see Section E.235.
E.232.2. Changes

- Enforce restrictions in `plperl` using an opmask applied to the whole interpreter, instead of using `Safe.pm` (Tim Bunce, Andrew Dunstan)

  Recent developments have convinced us that `Safe.pm` is too insecure to rely on for making `plperl` trustable. This change removes use of `Safe.pm` altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include that it is now possible to use Perl's `strict` pragma in a natural way in `plperl`, and that Perl's `$a` and `$b` variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

- Prevent PL/Tcl from executing untrustworthy code from `pltcl_modules` (Tom)

  PL/Tcl's feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless `pltcl_modules` is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted “normal” Tcl interpreter unless we are really going to execute a `pltclu` function. (CVE-2010-1170)

- Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

  Previously, if an unprivileged user ran `ALTER USER ... RESET ALL` for himself, or `ALTER DATABASE ... RESET ALL` for a database he owns, this would remove all special parameter settings for the user or database, even ones that are only supposed to be changeable by a superuser. Now, the `ALTER` will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a `CONTEXT` addition would be made to log entries (Tom)

  In some cases the context-printing function would fail because the current transaction had already been rolled back when it came time to print a log message.

- Update PL/Perl's `ppport.h` for modern Perl versions (Andrew)

- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)

- Prevent infinite recursion in `psql` when expanding a variable that refers to itself (Tom)

- Ensure that `contrib/pgstatstuple` functions respond to cancel interrupts promptly (Tatsuhito Kasahara)

- Make server startup deal properly with the case that `shmget()` returns `EINVAL` for an existing shared memory segment (Tom)

  This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entirely-misleading startup failure complaining that the shared memory request size was too large.

- Update time zone data files to `tzdata` release 2010j for DST law changes in Argentina, Australian Antarctic, Bangladesh, Mexico, Morocco, Pakistan, Palestine, Russia, Syria, Tunisia; also historical corrections for Taiwan.

E.233. Release 8.1.20

Release date: 2010-03-15

This release contains a variety of fixes from 8.1.19. For information about new features in the 8.1 major release, see Section E.253.
E.233.1. Migration to Version 8.1.20

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.18, see Section E.235.

E.233.2. Changes

• Add new configuration parameter ssl_renegotiation_limit to control how often we do session key renegotiation for an SSL connection (Magnus)

  This can be set to zero to disable renegotiation completely, which may be required if a broken SSL library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that cause renegotiation attempts to fail.

• Fix possible crashes when trying to recover from a failure in subtransaction start (Tom)

• Fix server memory leak associated with use of savepoints and a client encoding different from server's encoding (Tom)

• Make substring() for bit types treat any negative length as meaning “all the rest of the string” (Tom)

  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

• Fix integer-to-bit-string conversions to handle the first fractional byte correctly when the output bit width is wider than the given integer by something other than a multiple of 8 bits (Tom)

• Fix some cases of pathologically slow regular expression matching (Tom)

• Fix the STOP WAL LOCATION entry in backup history files to report the next WAL segment's name when the end location is exactly at a segment boundary (Itagaki Takahiro)

• Fix some more cases of temporary-file leakage (Heikki)

  This corrects a problem introduced in the previous minor release. One case that failed is when a plpgsql function returning set is called within another function's exception handler.

• When reading pg_hba.conf and related files, do not treat @something as a file inclusion request if the @ appears inside quote marks; also, never treat @ by itself as a file inclusion request (Tom)

  This prevents erratic behavior if a role or database name starts with @. If you need to include a file whose path name contains spaces, you can still do so, but you must write @"/path to/file" rather than putting the quotes around the whole construct.

• Prevent infinite loop on some platforms if a directory is named as an inclusion target in pg_hba.conf and related files (Tom)

• Fix psql's numericlocale option to not format strings it shouldn't in latex and troff output formats (Heikki)

• Fix plpgsql failure in one case where a composite column is set to NULL (Tom)

• Add volatile markings in PL/Python to avoid possible compiler-specific misbehavior (Zdenek Kotala)

• Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)

  The only known symptom of this oversight is that the Tcl clock command misbehaves if using Tcl 8.5 or later.
• Prevent crash in contrib/dblink when too many key columns are specified to a dblink_build_sql_* function (Rushabh Lathia, Joe Conway)

• Fix assorted crashes in contrib/xml2 caused by sloppy memory management (Tom)

• Update time zone data files to tzdata release 2010e for DST law changes in Bangladesh, Chile, Fiji, Mexico, Paraguay, Samoa.

E.234. Release 8.1.19

Release date: 2009-12-14

This release contains a variety of fixes from 8.1.18. For information about new features in the 8.1 major release, see Section E.253.

E.234.1. Migration to Version 8.1.19

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.18, see Section E.235.

E.234.2. Changes

• Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)

  This change prevents allegedly-immutable index functions from possibly subverting a superuser's session (CVE-2009-4136).

• Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)

  This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).

• Fix possible crash during backend-startup-time cache initialization (Tom)

• Prevent signals from interrupting VACUUM at unsafe times (Alvaro)

  This fix prevents a PANIC if a VACUUM FULL is canceled after it's already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

• Fix possible crash due to integer overflow in hash table size calculation (Tom)

  This could occur with extremely large planner estimates for the size of a hashjoin’s result.

• Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)

• Ensure that shared tuple-level locks held by prepared transactions are not ignored (Heikki)

• Fix premature drop of temporary files used for a cursor that is accessed within a subtransaction (Heikki)

• Fix PAM password processing to be more robust (Tom)

  The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.

• Fix processing of ownership dependencies during CREATE OR REPLACE FUNCTION (Tom)
• Ensure that Perl arrays are properly converted to PostgreSQL arrays when returned by a set-returning PL/Perl function (Andrew Dunstan, Abhijit Menon-Sen)

This worked correctly already for non-set-returning functions.

• Fix rare crash in exception processing in PL/Python (Peter)

• Ensure psql's flex module is compiled with the correct system header definitions (Tom)

This fixes build failures on platforms where --enable-largefile causes incompatible changes in the generated code.

• Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

• Update time zone data files to tzdata release 2009s for DST law changes in Antarctica, Argentina, Bangladesh, Fiji, Novokuznetsk, Pakistan, Palestine, Samoa, Syria; also historical corrections for Hong Kong.

E.235. Release 8.1.18

Release date: 2009-09-09

This release contains a variety of fixes from 8.1.17. For information about new features in the 8.1 major release, see Section E.253.

E.235.1. Migration to Version 8.1.18

A dump/restore is not required for those running 8.1.X. However, if you have any hash indexes on interval columns, you must REINDEX them after updating to 8.1.18. Also, if you are upgrading from a version earlier than 8.1.15, see Section E.238.

E.235.2. Changes

• Disallow RESET ROLE and RESET SESSION AUTHORIZATION inside security-definer functions (Tom, Heikki)

This covers a case that was missed in the previous patch that disallowed SET ROLE and SET SESSION AUTHORIZATION inside security-definer functions. (See CVE-2007-6600)

• Fix handling of sub-SELECTs appearing in the arguments of an outer-level aggregate function (Tom)

• Fix hash calculation for data type interval (Tom)

This corrects wrong results for hash joins on interval values. It also changes the contents of hash indexes on interval columns. If you have any such indexes, you must REINDEX them after updating.

• Treat to_char(..., 'TH') as an uppercase ordinal suffix with 'HH'/'HH12' (Heikki)

It was previously handled as 'th' (lowercase).

• Fix overflow for INTERVAL 'x ms' when x is more than 2 million and integer datetimes are in use (Alex Hunsaker)

• Fix calculation of distance between a point and a line segment (Tom)

This led to incorrect results from a number of geometric operators.
• Fix money data type to work in locales where currency amounts have no fractional digits, e.g. Japan (Itagaki Takahiro)

• Properly round datetime input like 00:12:57.9999999999999999999999999999 (Tom)

• Fix poor choice of page split point in GiST R-tree operator classes (Teodor)

• Fix portability issues in plperl initialization (Andrew Dunstan)

• Fix pg_ctl to not go into an infinite loop ifpostgresql.conf is empty (Jeff Davis)

• Fix contrib/xml2's xslt_process() to properly handle the maximum number of parameters (twenty) (Tom)

• Improve robustness of libpq's code to recover from errors during COPY FROM STDIN (Tom)

• Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)

• Update time zone data files to tzdata release 2009l for DST law changes in Bangladesh, Egypt, Jordan, Pakistan, Argentina/San_Luis, Cuba, Jordan (historical correction only), Mauritius, Morocco, Palestine, Syria, Tunisia.

E.236. Release 8.1.17

Release date: 2009-03-16

This release contains a variety of fixes from 8.1.16. For information about new features in the 8.1 major release, see Section E.253.

E.236.1. Migration to Version 8.1.17

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.15, see Section E.238.

E.236.2. Changes

• Prevent error recursion crashes when encoding conversion fails (Tom)

  This change extends fixes made in the last two minor releases for related failure scenarios. The previous fixes were narrowly tailored for the original problem reports, but we have now recognized that any error thrown by an encoding conversion function could potentially lead to infinite recursion while trying to report the error. The solution therefore is to disable translation and encoding conversion and report the plain-ASCII form of any error message, if we find we have gotten into a recursive error reporting situation. (CVE-2009-0922)

• Disallow CREATE CONVERSION with the wrong encodings for the specified conversion function (Heikki)

  This prevents one possible scenario for encoding conversion failure. The previous change is a backstop to guard against other kinds of failures in the same area.

• Fix core dump when to_char() is given format codes that are inappropriate for the type of the data argument (Tom)

• Fix decompilation of CASE WHEN with an implicit coercion (Tom)

  This mistake could lead to Assert failures in an Assert-enabled build, or an “unexpected CASE WHEN clause” error message in other cases, when trying to examine or dump a view.
• Fix possible misassignment of the owner of a TOAST table's rowtype (Tom)

If `CLUSTER` or a rewriting variant of `ALTER TABLE` were executed by someone other than the table owner, the `pg_type` entry for the table's TOAST table would end up marked as owned by that someone. This caused no immediate problems, since the permissions on the TOAST rowtype aren't examined by any ordinary database operation. However, it could lead to unexpected failures if one later tried to drop the role that issued the command (in 8.1 or 8.2), or "owner of data type appears to be invalid" warnings from `pg_dump` after having done so (in 8.3).

• Clean up PL/pgSQL error status variables fully at block exit (Ashesh Vashi and Dave Page)

This is not a problem for PL/pgSQL itself, but the omission could cause the PL/pgSQL Debugger to crash while examining the state of a function.

• Add **MUST** (Mauritius Island Summer Time) to the default list of known timezone abbreviations (Xavier Bugaud)

### E.237. Release 8.1.16

**Release date:** 2009-02-02

This release contains a variety of fixes from 8.1.15. For information about new features in the 8.1 major release, see Section E.253.

#### E.237.1. Migration to Version 8.1.16

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.15, see Section E.238.

#### E.237.2. Changes

• Fix crash in autovacuum (Alvaro)

  The crash occurs only after vacuuming a whole database for anti-transaction-wraparound purposes, which means that it occurs infrequently and is hard to track down.

• Improve handling of URLs in `headline()` function (Teodor)

• Improve handling of overlength headlines in `headline()` function (Teodor)

• Prevent possible Assert failure or misconversion if an encoding conversion is created with the wrong conversion function for the specified pair of encodings (Tom, Heikki)

• Avoid unnecessary locking of small tables in `VACUUM` (Heikki)

• Ensure that the contents of a holdable cursor don't depend on the contents of TOAST tables (Tom)

  Previously, large field values in a cursor result might be represented as TOAST pointers, which would fail if the referenced table got dropped before the cursor is read, or if the large value is deleted and then vacuumed away. This cannot happen with an ordinary cursor, but it could with a cursor that is held past its creating transaction.

• Fix uninitialized variables in `contrib/tsearch2's get_covers()` function (Teodor)

• Fix configure script to properly report failure when unable to obtain linkage information for PL/Perl (Andrew)

• Make all documentation reference `pgsql-bugs` and/or `pgsql-hackers` as appropriate, instead of the now-decommissioned `pgsql-ports` and `pgsql-patches` mailing lists (Tom)
• Update time zone data files to tzdata release 2009a (for Kathmandu and historical DST corrections in Switzerland, Cuba)

E.238. Release 8.1.15

Release date: 2008-11-03

This release contains a variety of fixes from 8.1.14. For information about new features in the 8.1 major release, see Section E.253.

E.238.1. Migration to Version 8.1.15

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251. Also, if you were running a previous 8.1.X release, it is recommended to REINDEX all GiST indexes after the upgrade.

E.238.2. Changes

• Fix GiST index corruption due to marking the wrong index entry “dead” after a deletion (Teodor)

This would result in index searches failing to find rows they should have found. Corrupted indexes can be fixed with REINDEX.

• Fix backend crash when the client encoding cannot represent a localized error message (Tom)

We have addressed similar issues before, but it would still fail if the “character has no equivalent” message itself couldn’t be converted. The fix is to disable localization and send the plain ASCII error message when we detect such a situation.

• Fix possible crash when deeply nested functions are invoked from a trigger (Tom)

• Fix mis-expansion of rule queries when a sub-SELECT appears in a function call in FROM, a multi-row VALUES list, or a RETURNING list (Tom)

The usual symptom of this problem is an “unrecognized node type” error.

• Ensure an error is reported when a newly-defined PL/pgSQL trigger function is invoked as a normal function (Tom)

• Prevent possible collision of relfilenode numbers when moving a table to another tablespace with ALTER SET TABLESPACE (Heikki)

The command tried to re-use the existing filename, instead of picking one that is known unused in the destination directory.

• Fix incorrect tssearch2 headline generation when single query item matches first word of text (Sushant Sinha)

• Fix improper display of fractional seconds in interval values when using a non-ISO datestyle in an --enable-integer-datetime build (Ron Mayer)

• Ensure SPI_getvalue and SPI_getbinval behave correctly when the passed tuple and tuple descriptor have different numbers of columns (Tom)

This situation is normal when a table has had columns added or removed, but these two functions didn’t handle it properly. The only likely consequence is an incorrect error indication.

• Fix ecpg's parsing of CREATE ROLE (Michael)
• Fix recent breakage of `pg_ctl restart` (Tom)

• Update time zone data files to `tzdata` release 2008i (for DST law changes in Argentina, Brazil, Mauritius, Syria)

E.239. Release 8.1.14

Release date: 2008-09-22

This release contains a variety of fixes from 8.1.13. For information about new features in the 8.1 major release, see Section E.253.

E.239.1. Migration to Version 8.1.14

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.239.2. Changes

• Widen local lock counters from 32 to 64 bits (Tom)

  This responds to reports that the counters could overflow in sufficiently long transactions, leading to unexpected “lock is already held” errors.

• Fix possible duplicate output of tuples during a GiST index scan (Teodor)

• Add checks in executor startup to ensure that the tuples produced by an `INSERT` or `UPDATE` will match the target table’s current rowtype (Tom)

  `ALTER COLUMN TYPE`, followed by re-use of a previously cached plan, could produce this type of situation. The check protects against data corruption and/or crashes that could ensue.

• Fix `AT TIME ZONE` to first try to interpret its timezone argument as a timezone abbreviation, and only try it as a full timezone name if that fails, rather than the other way around as formerly (Tom)

  The timestamp input functions have always resolved ambiguous zone names in this order. Making `AT TIME ZONE` do so as well improves consistency, and fixes a compatibility bug introduced in 8.1: in ambiguous cases we now behave the same as 8.0 and before did, since in the older versions `AT TIME ZONE` accepted *only* abbreviations.

• Fix datetime input functions to correctly detect integer overflow when running on a 64-bit platform (Tom)

• Improve performance of writing very long log messages to syslog (Tom)

• Fix bug in backwards scanning of a cursor on a `SELECT DISTINCT ON` query (Tom)

• Fix planner bug with nested sub-select expressions (Tom)

  If the outer sub-select has no direct dependency on the parent query, but the inner one does, the outer value might not get recalculated for new parent query rows.

• Fix planner to estimate that `GROUP BY` expressions yielding boolean results always result in two groups, regardless of the expressions' contents (Tom)

  This is very substantially more accurate than the regular `GROUP BY` estimate for certain boolean tests like `col IS NULL`.

• Fix PL/pgSQL to not fail when a `FOR` loop's target variable is a record containing composite-type fields (Tom)
• Fix PL/Tcl to behave correctly with Tcl 8.5, and to be more careful about the encoding of data sent to or from Tcl (Tom)

• Fix PL/Python to work with Python 2.5

    This is a back-port of fixes made during the 8.2 development cycle.

• Improve pg_dump and pg_restore's error reporting after failure to send a SQL command (Tom)

• Fix pg_ctl to properly preserve postmaster command-line arguments across a restart (Bruce)

• Update time zone data files to tzdata release 2008f (for DST law changes in Argentina, Bahamas, Brazil, Mauritius, Morocco, Pakistan, Palestine, and Paraguay)

E.240. Release 8.1.13

Release date: 2008-06-12

This release contains one serious and one minor bug fix over 8.1.12. For information about new features in the 8.1 major release, see Section E.253.

E.240.1. Migration to Version 8.1.13

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.240.2. Changes

• Make \texttt{pg\_get\_ruledef()} parenthesize negative constants (Tom)

    Before this fix, a negative constant in a view or rule might be dumped as, say, \texttt{-42::integer}, which is subtly incorrect: it should be \texttt{(-42)::integer} due to operator precedence rules. Usually this would make little difference, but it could interact with another recent patch to cause PostgreSQL to reject what had been a valid \texttt{SELECT DISTINCT} view query. Since this could result in pg_dump output failing to reload, it is being treated as a high-priority fix. The only released versions in which dump output is actually incorrect are 8.3.1 and 8.2.7.

• Make \texttt{ALTER AGGREGATE ... OWNER TO} update \texttt{pg\_shdepend} (Tom)

    This oversight could lead to problems if the aggregate was later involved in a \texttt{DROP OWNED} or \texttt{REASSIGN OWNED} operation.


Release date: never released

This release contains a variety of fixes from 8.1.11. For information about new features in the 8.1 major release, see Section E.253.

E.241.1. Migration to Version 8.1.12

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.241.2. Changes

• Fix \texttt{ALTER TABLE ADD COLUMN ... PRIMARY KEY} so that the new column is correctly checked to see if it's been initialized to all non-nulls (Brendan Jurd)
Previous versions neglected to check this requirement at all.

- Fix possible `CREATE TABLE` failure when inheriting the “same” constraint from multiple parent relations that inherited that constraint from a common ancestor (Tom)

- Fix conversions between ISO-8859-5 and other encodings to handle Cyrillic “Yo” characters (e and E with two dots) (Sergey Burladyan)

- Fix a few datatype input functions that were allowing unused bytes in their results to contain uninitialized, unpredictable values (Tom)

  This could lead to failures in which two apparently identical literal values were not seen as equal, resulting in the parser complaining about unmatched `ORDER BY` and `DISTINCT` expressions.

- Fix a corner case in regular-expression substring matching (`substring(string from pattern)`) (Tom)

  The problem occurs when there is a match to the pattern overall but the user has specified a parenthesized subexpression and that subexpression hasn’t got a match. An example is `substring('foo' from 'foo(bar)?')`. This should return NULL, since (bar) isn’t matched, but it was mistakenly returning the whole-pattern match instead (ie, foo).

- Fix time zone data files to tzdata release 2008c (for DST law changes in Morocco, Iraq, Choibalsan, Pakistan, Syria, Cuba, Argentina/San_Luis, and Chile)

- Fix incorrect result from ecpg’s `PGTYPEStimestamp_sub` function (Michael)

- Fix core dump in contrib/xml2’s `xpath_table()` function when the input query returns a NULL value (Tom)

- Fix contrib/xml2’s makefile to not override `CFLAGS` (Tom)

- Fix `DatumGetBool` macro to not fail with gcc 4.3 (Tom)

  This problem affects “old style” (V0) C functions that return boolean. The fix is already in 8.3, but the need to back-patch it was not realized at the time.

- Fix longstanding `LISTEN`/`NOTIFY` race condition (Tom)

  In rare cases a session that had just executed a `LISTEN` might not get a notification, even though one would be expected because the concurrent transaction executing `NOTIFY` was observed to commit later.

  A side effect of the fix is that a transaction that has executed a not-yet-committed `LISTEN` command will not see any row in `pg_listener` for the `LISTEN`, should it choose to look; formerly it would have. This behavior was never documented one way or the other, but it is possible that some applications depend on the old behavior.

- Disallow `LISTEN` and `UNLISTEN` within a prepared transaction (Tom)

  This was formerly allowed but trying to do it had various unpleasant consequences, notably that the originating backend could not exit as long as an `UNLISTEN` remained uncommitted.

- Fix rare crash when an error occurs during a query using a hash index (Heikki)

- Fix input of datetime values for February 29 in years BC (Tom)

  The former coding was mistaken about which years were leap years.

- Fix “unrecognized node type” error in some variants of `ALTER OWNER` (Tom)
• Fix pg_ctl to correctly extract the postmaster's port number from command-line options (Itagaki Takahiro, Tom)

Previously, `pg_ctl start -w` could try to contact the postmaster on the wrong port, leading to bogus reports of startup failure.

• Use `-fwrapv` to defend against possible misoptimization in recent gcc versions (Tom)

This is known to be necessary when building PostgreSQL with gcc 4.3 or later.

• Fix display of constant expressions in `ORDER BY` and `GROUP BY` (Tom)

An explicitly casted constant would be shown incorrectly. This could for example lead to corruption of a view definition during dump and reload.

• Fix libpq to handle NOTICE messages correctly during COPY OUT (Tom)

This failure has only been observed to occur when a user-defined datatype's output routine issues a NOTICE, but there is no guarantee it couldn't happen due to other causes.

E.242. Release 8.1.11

Release date: 2008-01-07

This release contains a variety of fixes from 8.1.10, including fixes for significant security issues. For information about new features in the 8.1 major release, see Section E.253.

This is the last 8.1.X release for which the PostgreSQL community will produce binary packages for Windows. Windows users are encouraged to move to 8.2.X or later, since there are Windows-specific fixes in 8.2.X that are impractical to back-port. 8.1.X will continue to be supported on other platforms.

E.242.1. Migration to Version 8.1.11

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.242.2. Changes

• Prevent functions in indexes from executing with the privileges of the user running VACUUM, ANALYZE, etc (Tom)

Functions used in index expressions and partial-index predicates are evaluated whenever a new table entry is made. It has long been understood that this poses a risk of trojan-horse code execution if one modifies a table owned by an untrustworthy user. (Note that triggers, defaults, check constraints, etc. pose the same type of risk.) But functions in indexes pose extra danger because they will be executed by routine maintenance operations such as VACUUM FULL, which are commonly performed automatically under a superuser account. For example, a nefarious user can execute code with superuser privileges by setting up a trojan-horse index definition and waiting for the next routine vacuum. The fix arranges for standard maintenance operations (including VACUUM, ANALYZE, REINDEX, and CLUSTER) to execute as the table owner rather than the calling user, using the same privilege-switching mechanism already used for SECURITY DEFINER functions. To prevent bypassing this security measure, execution of SET SESSION AUTHORIZATION and SET ROLE is now forbidden within a SECURITY DEFINER context. (CVE-2007-6600)

• Repair assorted bugs in the regular-expression package (Tom, Will Drewry)

Suitably crafted regular-expression patterns could cause crashes, infinite or near-infinite looping, and/or massive memory consumption, all of which pose denial-of-service hazards for applications.

• Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

  The fix that appeared for this in 8.1.10 was incomplete, as it plugged the hole for only some dblink functions. (CVE-2007-6601, CVE-2007-3278)

• Update time zone data files to tzdata release 2007k (in particular, recent Argentina changes) (Tom)

• Improve planner’s handling of LIKE/regex estimation in non-C locales (Tom)

• Fix planner failure in some cases of WHERE false AND var IN (SELECT ...) (Tom)

• Preserve the tablespace of indexes that are rebuilt by ALTER TABLE ... ALTER COLUMN TYPE (Tom)

• Make archive recovery always start a new WAL timeline, rather than only when a recovery stop time was used (Simon)

  This avoids a corner-case risk of trying to overwrite an existing archived copy of the last WAL segment, and seems simpler and cleaner than the original definition.

• Make VACUUM not use all of maintenance_work_mem when the table is too small for it to be useful (Alvaro)

• Fix potential crash in translate() when using a multibyte database encoding (Tom)

• Fix overflow in extract (epoch from interval) for intervals exceeding 68 years (Tom)

• Fix PL/Perl to not fail when a UTF-8 regular expression is used in a trusted function (Andrew)

• Fix PL/Perl to cope when platform's Perl defines type bool as int rather than char (Tom)

  While this could theoretically happen anywhere, no standard build of Perl did things this way ... until macOS 10.5.

• Fix PL/Python to not crash on long exception messages (Alvaro)

• Fix pg_dump to correctly handle inheritance child tables that have default expressions different from their parent’s (Tom)

• Fix libpq crash when PGPASSFILE refers to a file that is not a plain file (Martin Pitt)

• ecpg parser fixes (Michael)

• Make contrib/pgcrypto defend against OpenSSL libraries that fail on keys longer than 128 bits; which is the case at least on some Solaris versions (Marko Kreen)

• Make contrib/tablefunc's crosstab() handle NULL rowid as a category in its own right, rather than crashing (Joe)

• Fix tsvector and tsquery output routines to escape backslashes correctly (Teodor, Bruce)

• Fix crash of to_tsvector() on huge input strings (Teodor)

• Require a specific version of Autoconf to be used when re-generating the configure script (Peter)

  This affects developers and packagers only. The change was made to prevent accidental use of untested combinations of Autoconf and PostgreSQL versions. You can remove the version check

3096
if you really want to use a different Autoconf version, but it's your responsibility whether the result works or not.

**E.243. Release 8.1.10**

**Release date:** 2007-09-17

This release contains a variety of fixes from 8.1.9. For information about new features in the 8.1 major release, see Section E.253.

**E.243.1. Migration to Version 8.1.10**

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

**E.243.2. Changes**

- Prevent index corruption when a transaction inserts rows and then aborts close to the end of a concurrent VACUUM on the same table (Tom)
- Make CREATE DOMAIN ... DEFAULT NULL work properly (Tom)
- Allow the interval data type to accept input consisting only of milliseconds or microseconds (Neil)
- Speed up rtree index insertion (Teodor)
- Fix excessive logging of SSL error messages (Tom)
- Fix logging so that log messages are never interleaved when using the syslogger process (Andrew)
- Fix crash when log_min_error_statement logging runs out of memory (Tom)
- Fix incorrect handling of some foreign-key corner cases (Tom)
- Prevent REINDEX and CLUSTER from failing due to attempting to process temporary tables of other sessions (Alvaro)
- Update the time zone database rules, particularly New Zealand's upcoming changes (Tom)
- Windows socket improvements (Magnus)
- Suppress timezone name (%Z) in log timestamps on Windows because of possible encoding mismatches (Tom)
- Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

**E.244. Release 8.1.9**

**Release date:** 2007-04-23

This release contains a variety of fixes from 8.1.8, including a security fix. For information about new features in the 8.1 major release, see Section E.253.

**E.244.1. Migration to Version 8.1.9**

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.
E.244.2. Changes

- Support explicit placement of the temporary-table schema within search_path, and disable searching it for functions and operators (Tom)
  
  This is needed to allow a security-definer function to set a truly secure value of search_path. Without it, an unprivileged SQL user can use temporary objects to execute code with the privileges of the security-definer function (CVE-2007-2138). See CREATE FUNCTION for more information.

- /contrib/tsearch2 crash fixes (Teodor)

- Require COMMIT PREPARED to be executed in the same database as the transaction was prepared in (Heikki)

- Fix potential-data-corruption bug in how VACUUM FULL handles UPDATE chains (Tom, Pavan Deolasee)

- Planner fixes, including improving outer join and bitmap scan selection logic (Tom)

- Fix PANIC during enlargement of a hash index (bug introduced in 8.1.6) (Tom)

- Fix POSIX-style timezone specs to follow new USA DST rules (Tom)

E.245. Release 8.1.8

Release date: 2007-02-07

This release contains one fix from 8.1.7. For information about new features in the 8.1 major release, see Section E.253.

E.245.1. Migration to Version 8.1.8

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.245.2. Changes

- Remove overly-restrictive check for type length in constraints and functional indexes (Tom)

E.246. Release 8.1.7

Release date: 2007-02-05

This release contains a variety of fixes from 8.1.6, including a security fix. For information about new features in the 8.1 major release, see Section E.253.

E.246.1. Migration to Version 8.1.7

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.246.2. Changes

- Remove security vulnerabilities that allowed connected users to read backend memory (Tom)

  The vulnerabilities involve suppressing the normal check that a SQL function returns the data type it's declared to, and changing the data type of a table column (CVE-2007-0555, CVE-2007-0556).
These errors can easily be exploited to cause a backend crash, and in principle might be used to read database content that the user should not be able to access.

- Fix rare bug wherein btree index page splits could fail due to choosing an infeasible split point (Heikki Linnakangas)
- Improve VACUUM performance for databases with many tables (Tom)
- Fix autovacuum to avoid leaving non-permanent transaction IDs in non-connectable databases (Alvaro)

This bug affects the 8.1 branch only.

- Fix for rare Assert() crash triggered by UNION (Tom)
- Tighten security of multi-byte character processing for UTF8 sequences over three bytes long (Tom)
- Fix bogus “permission denied” failures occurring on Windows due to attempts to fsync already-deleted files (Magnus, Tom)
- Fix possible crashes when an already-in-use PL/pgSQL function is updated (Tom)

E.247. Release 8.1.6

Release date: 2007-01-08

This release contains a variety of fixes from 8.1.5. For information about new features in the 8.1 major release, see Section E.253.

E.247.1. Migration to Version 8.1.6

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.247.2. Changes

- Improve handling of getaddrinfo() on AIX (Tom)
  This fixes a problem with starting the statistics collector, among other things.
- Fix pg_restore to handle a tar-format backup that contains large objects (blobs) with comments (Tom)
- Fix “failed to re-find parent key” errors in VACUUM (Tom)
- Clean out pg_internal.init cache files during server restart (Simon)
  This avoids a hazard that the cache files might contain stale data after PITR recovery.
- Fix race condition for truncation of a large relation across a gigabyte boundary by VACUUM (Tom)
- Fix bug causing needless deadlock errors on row-level locks (Tom)
- Fix bugs affecting multi-gigabyte hash indexes (Tom)
- Fix possible deadlock in Windows signal handling (Teodor)
- Fix error when constructing an ARRAY[] made up of multiple empty elements (Tom)
- Fix ecpg memory leak during connection (Michael)
- Fix for macOS (Darwin) compilation (Tom)
• `to_number()` and `to_char(numeric)` are now STABLE, not IMMUTABLE, for new `initdb` installs (Tom)
  
  This is because `lc_numeric` can potentially change the output of these functions.

• Improve index usage of regular expressions that use parentheses (Tom)
  
  This improves `psql` \d performance also.

• Update timezone database
  
  This affects Australian and Canadian daylight-savings rules in particular.

E.248. Release 8.1.5

Release date: 2006-10-16

This release contains a variety of fixes from 8.1.4. For information about new features in the 8.1 major release, see Section E.253.

E.248.1. Migration to Version 8.1.5

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.248.2. Changes

• Disallow aggregate functions in `UPDATE` commands, except within sub-`SELECT`s (Tom)

  The behavior of such an aggregate was unpredictable, and in 8.1.X could cause a crash, so it has been disabled. The SQL standard does not allow this either.

• Fix core dump when an untyped literal is taken as `ANYARRAY`

• Fix core dump in duration logging for extended query protocol when a `COMMIT` or `ROLLBACK` is executed

• Fix mishandling of `AFTER` triggers when query contains a SQL function returning multiple rows (Tom)

• Fix `ALTER TABLE ... TYPE` to recheck `NOT NULL` for `USING` clause (Tom)

• Fix `string_to_array()` to handle overlapping matches for the separator string

  For example, `string_to_array('123xx456xxx789', 'xx').`

• Fix `to_timestamp()` for AM/PM formats (Bruce)

• Fix `autovacuum`'s calculation that decides whether `ANALYZE` is needed (Alvaro)

• Fix corner cases in pattern matching for `psql`'s \d commands

• Fix index-corrupting bugs in `/contrib/ltree` (Teodor)

• Numerous robustness fixes in `ecpg` (Joachim Wieland)

• Fix backslash escaping in `/contrib/dbmirror`

• Minor fixes in `/contrib/dblink` and `/contrib/tsearch2`

• Efficiency improvements in hash tables and bitmap index scans (Tom)
• Fix instability of statistics collection on Windows (Tom, Andrew)

• Fix statement_timeout to use the proper units on Win32 (Bruce)

  In previous Win32 8.1.X versions, the delay was off by a factor of 100.

• Fixes for MSVC and Borland C++ compilers (Hiroshi Saito)

• Fixes for AIX and Intel compilers (Tom)

• Fix rare bug in continuous archiving (Tom)

E.249. Release 8.1.4

Release date: 2006-05-23

This release contains a variety of fixes from 8.1.3, including patches for extremely serious security issues. For information about new features in the 8.1 major release, see Section E.253.

E.249.1. Migration to Version 8.1.4

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

Full security against the SQL-injection attacks described in CVE-2006-2313 and CVE-2006-2314 might require changes in application code. If you have applications that embed untrustworthy strings into SQL commands, you should examine them as soon as possible to ensure that they are using recommended escaping techniques. In most cases, applications should be using subroutines provided by libraries or drivers (such as libpq's PQescapeStringConn()) to perform string escaping, rather than relying on ad hoc code to do it.

E.249.2. Changes

• Change the server to reject invalidly-encoded multibyte characters in all cases (Tatsuo, Tom)

  While PostgreSQL has been moving in this direction for some time, the checks are now applied uniformly to all encodings and all textual input, and are now always errors not merely warnings. This change defends against SQL-injection attacks of the type described in CVE-2006-2313.

• Reject unsafe uses of \' in string literals

  As a server-side defense against SQL-injection attacks of the type described in CVE-2006-2314, the server now only accepts ' and not \' as a representation of ASCII single quote in SQL string literals. By default, \' is rejected only when client_encoding is set to a client-only encoding (SJIS, BIG5, GBK, GB18030, or UHC), which is the scenario in which SQL injection is possible. A new configuration parameter backslash_quote is available to adjust this behavior when needed. Note that full security against CVE-2006-2314 might require client-side changes; the purpose of backslash_quote is in part to make it obvious that insecure clients are insecure.

• Modify libpq's string-escaping routines to be aware of encoding considerations and standard_conforming_strings

  This fixes libpq-using applications for the security issues described in CVE-2006-2313 and CVE-2006-2314, and also future-proofs them against the planned changeover to SQL-standard string literal syntax. Applications that use multiple PostgreSQL connections concurrently should migrate to PQescapeStringConn() and PQescapeByteaConn() to ensure that escaping is done correctly for the settings in use in each database connection. Applications that do string escaping "by hand" should be modified to rely on library routines instead.
• Fix weak key selection in pgcrypto (Marko Kreen)

Errors in fortuna PRNG reseeding logic could cause a predictable session key to be selected by pgp_sym_encrypt() in some cases. This only affects non-OpenSSL-using builds.

• Fix some incorrect encoding conversion functions

win1251_to_iso, win866_to_iso, euc_tw_to_big5, euc_tw_to_mic, mic_to_euc_tw were all broken to varying extents.

• Clean up stray remaining uses of \' in strings (Bruce, Jan)

• Make autovacuum visible in pg_stat_activity (Alvaro)

• Disable full_page_writes (Tom)

In certain cases, having full_page_writes off would cause crash recovery to fail. A proper fix will appear in 8.2; for now it's just disabled.

• Various planner fixes, particularly for bitmap index scans and MIN/MAX optimization (Tom)

• Fix incorrect optimization in merge join (Tom)

Outer joins could sometimes emit multiple copies of unmatched rows.

• Fix crash from using and modifying a plpgsql function in the same transaction

• Fix WAL replay for case where a B-Tree index has been truncated

• Fix SIMILAR TO for patterns involving | (Tom)

• Fix SELECT INTO and CREATE TABLE AS to create tables in the default tablespace, not the base directory (Kris Jurka)

• Fix server to use custom DH SSL parameters correctly (Michael Fuhr)

• Improve qsort performance (Dann Corbit)

Currently this code is only used on Solaris.

• Fix for OS/X Bonjour on x86 systems (Ashley Clark)

• Fix various minor memory leaks

• Fix problem with password prompting on some Win32 systems (Robert Kinberg)

• Improve pg_dump's handling of default values for domains

• Fix pg_dumpall to handle identically-named users and groups reasonably (only possible when dumping from a pre-8.1 server) (Tom)

The user and group will be merged into a single role with LOGIN permission. Formerly the merged role wouldn't have LOGIN permission, making it unusable as a user.

• Fix pg_restore -n to work as documented (Tom)

E.250. Release 8.1.3

Release date: 2006-02-14

This release contains a variety of fixes from 8.1.2, including one very serious security issue. For information about new features in the 8.1 major release, see Section E.253.
E.250.1. Migration to Version 8.1.3

A dump/restore is not required for those running 8.1.X. However, if you are upgrading from a version earlier than 8.1.2, see Section E.251.

E.250.2. Changes

- Fix bug that allowed any logged-in user to `SET ROLE` to any other database user id (CVE-2006-0553)

  Due to inadequacy of validation checking, a user could exploit the special case that `SET ROLE` normally uses to restore the previous role setting after an error. This allowed ordinary users to acquire superuser status, for example. The escalation-of-privilege risk exists only in 8.1.0-8.1.2. However, in all releases back to 7.3 there is a related bug in `SET SESSION AUTHORIZATION` that allows unprivileged users to crash the server, if it has been compiled with Asserts enabled (which is not the default). Thanks to Akio Ishida for reporting this problem.

- Fix bug with row visibility logic in self-inserted rows (Tom)

  Under rare circumstances a row inserted by the current command could be seen as already valid, when it should not be. Repairs bug created in 8.0.4, 7.4.9, and 7.3.11 releases.

- Fix race condition that could lead to “file already exists” errors during `pg_clog` and `pg_subtrans` file creation (Tom)

- Fix cases that could lead to crashes if a cache-invalidation message arrives at just the wrong time (Tom)

- Properly check `DOMAIN` constraints for `UNKNOWN` parameters in prepared statements (Neil)

- Ensure `ALTER COLUMN TYPE` will process `FOREIGN KEY`, `UNIQUE`, and `PRIMARY KEY` constraints in the proper order (Nakano Yoshihisa)

- Fixes to allow restoring dumps that have cross-schema references to custom operators or operator classes (Tom)

- Allow `pg_restore` to continue properly after a `COPY` failure; formerly it tried to treat the remaining `COPY` data as SQL commands (Stephen Frost)

- Fix `pg_ctl unregister` crash when the data directory is not specified (Magnus)

- Fix `libpq PQprint` HTML tags (Christoph Zwerschke)

- Fix `ecpg` crash on AMD64 and PPC (Neil)

- Allow `SETOF` and `%TYPE` to be used together in function result type declarations

- Recover properly if error occurs during argument passing in PL/Python (Neil)

- Fix memory leak in `plperl_return_next` (Neil)

- Fix PL/Perl’s handling of locales on Win32 to match the backend (Andrew)

- Various optimizer fixes (Tom)

- Fix crash when `log_min_messages` is set to `DEBUG3` or above in `postgresql.conf` on Win32 (Bruce)

- Fix `pgxs -L` library path specification for Win32, Cygwin, macOS, AIX (Bruce)

- Check that `SID` is enabled while checking for Win32 admin privileges (Magnus)
• Properly reject out-of-range date inputs (Kris Jurka)
• Portability fix for testing presence of finite and isinf during configure (Tom)
• Improve speed of COPY IN via libpq, by avoiding a kernel call per data line (Alon Goldshuv)
• Improve speed of /contrib/tsearch2 index creation (Tom)

E.251. Release 8.1.2

Release date: 2006-01-09

This release contains a variety of fixes from 8.1.1. For information about new features in the 8.1 major release, see Section E.253.

E.251.1. Migration to Version 8.1.2

A dump/restore is not required for those running 8.1.X. However, you might need to REINDEX indexes on textual columns after updating, if you are affected by the locale or plperl issues described below.

E.251.2. Changes

• Fix Windows code so that postmaster will continue rather than exit if there is no more room in ShmemBackendArray (Magnus)

  The previous behavior could lead to a denial-of-service situation if too many connection requests arrive close together. This applies only to the Windows port.

• Fix bug introduced in 8.0 that could allow ReadBuffer to return an already-used page as new, potentially causing loss of recently-committed data (Tom)

• Fix for protocol-level Describe messages issued outside a transaction or in a failed transaction (Tom)

• Fix character string comparison for locales that consider different character combinations as equal, such as Hungarian (Tom)

  This might require REINDEX to fix existing indexes on textual columns.

• Set locale environment variables during postmaster startup to ensure that plperl won’t change the locale later

  This fixes a problem that occurred if the postmaster was started with environment variables specifying a different locale than what initdb had been told. Under these conditions, any use of plperl was likely to lead to corrupt indexes. You might need REINDEX to fix existing indexes on textual columns if this has happened to you.

• Allow more flexible relocation of installation directories (Tom)

  Previous releases supported relocation only if all installation directory paths were the same except for the last component.

• Prevent crashes caused by the use of ISO-8859-5 and ISO-8859-9 encodings (Tatsuo)

• Fix longstanding bug in strpos() and regular expression handling in certain rarely used Asian multi-byte character sets (Tatsuo)

• Fix bug where COPY CSV mode considered any \. to terminate the copy data

  The new code requires \. to appear alone on a line, as per documentation.
• Make COPY CSV mode quote a literal data value of "," to ensure it cannot be interpreted as the end-of-data marker (Bruce)

• Various fixes for functions returning RECORDs (Tom)

• Fix processing of postgresql.conf so a final line with no newline is processed properly (Tom)

• Fix bug in /contrib/pgcrypto gen_salt, which caused it not to use all available salt space for MD5 and XDES algorithms (Marko Kreen, Solar Designer)

Salts for Blowfish and standard DES are unaffected.

• Fix autovacuum crash when processing expression indexes

• Fix /contrib/dblink to throw an error, rather than crashing, when the number of columns specified is different from what's actually returned by the query (Joe)

E.252. Release 8.1.1

Release date: 2005-12-12

This release contains a variety of fixes from 8.1.0. For information about new features in the 8.1 major release, see Section E.253.

E.252.1. Migration to Version 8.1.1

A dump/restore is not required for those running 8.1.X.

E.252.2. Changes

• Fix incorrect optimizations of outer-join conditions (Tom)

• Fix problems with wrong reported column names in cases involving sub-selects flattened by the optimizer (Tom)

• Fix update failures in scenarios involving CHECK constraints, toasted columns, and indexes (Tom)

• Fix bgwriter problems after recovering from errors (Tom)

The background writer was found to leak buffer pins after write errors. While not fatal in itself, this might lead to mysterious blockages of later VACUUM commands.

• Prevent failure if client sends Bind protocol message when current transaction is already aborted

• /contrib/tsearch2 and /contrib/ltree fixes (Teodor)

• Fix problems with translated error messages in languages that require word reordering, such as Turkish; also problems with unexpected truncation of output strings and wrong display of the smallest possible bigint value (Andrew, Tom)

These problems only appeared on platforms that were using our port/snprintf.c code, which includes BSD variants if --enable-nls was given, and perhaps others. In addition, a different form of the translated-error-message problem could appear on Windows depending on which version of libintl was used.

• Re-allow AM/PM, HH, HH12, and D format specifiers for to_char(time) and to_char(interval).to_char(interval) should probably use HH24. (Bruce)

• AIX, HPUX, and MSVC compile fixes (Tom, Hiroshi Saito)
• Optimizer improvements (Tom)

• Retry file reads and writes after Windows NO_SYSTEM_RESOURCES error (Qingqing Zhou)

• Prevent autovacuum from crashing during ANALYZE of expression index (Alvaro)

• Fix problems with ON COMMIT DELETE ROWS temp tables

• Fix problems when a trigger alters the output of a SELECT DISTINCT query

• Add 8.1.0 release note item on how to migrate invalid UTF-8 byte sequences (Paul Lindner)

E.253. Release 8.1

Release date: 2005-11-08

E.253.1. Overview

Major changes in this release:

Improve concurrent access to the shared buffer cache (Tom)

Access to the shared buffer cache was identified as a significant scalability problem, particularly on multi-CPU systems. In this release, the way that locking is done in the buffer manager has been overhauled to reduce lock contention and improve scalability. The buffer manager has also been changed to use a “clock sweep” replacement policy.

Allow index scans to use an intermediate in-memory bitmap (Tom)

In previous releases, only a single index could be used to do lookups on a table. With this feature, if a query has WHERE tab.col1 = 4 and tab.col2 = 9, and there is no multicolurn index on col1 and col2, but there is an index on col1 and another on col2, it is possible to search both indexes and combine the results in memory, then do heap fetches for only the rows matching both the col1 and col2 restrictions. This is very useful in environments that have a lot of unstructured queries where it is impossible to create indexes that match all possible access conditions. Bitmap scans are useful even with a single index, as they reduce the amount of random access needed; a bitmap index scan is efficient for retrieving fairly large fractions of the complete table, whereas plain index scans are not.

Add two-phase commit (Heikki Linnakangas, Alvaro, Tom)

Two-phase commit allows transactions to be "prepared" on several computers, and once all computers have successfully prepared their transactions (none failed), all transactions can be committed. Even if a machine crashes after a prepare, the prepared transaction can be committed after the machine is restarted. New syntax includes PREPARE TRANSACTION and COMMIT/ROLLBACK PREPARED. A new system view pg_prepared_xacts has also been added.

Create a new role system that replaces users and groups (Stephen Frost)

Roles are a combination of users and groups. Like users, they can have login capability, and like groups, a role can have other roles as members. Roles basically remove the distinction between users and groups. For example, a role can:

• Have login capability (optionally)

• Own objects

• Hold access permissions for database objects

• Inherit permissions from other roles it is a member of
Once a user logs into a role, she obtains capabilities of the login role plus any inherited roles, and can use SET ROLE to switch to other roles she is a member of. This feature is a generalization of the SQL standard's concept of roles. This change also replaces pg_shadow and pg_group by new role-capable catalogs pg_authid and pg_auth_members. The old tables are redefined as read-only views on the new role tables.

Automatically use indexes for \texttt{MIN()} and \texttt{MAX()} (Tom)

In previous releases, the only way to use an index for \texttt{MIN()} or \texttt{MAX()} was to rewrite the query as \texttt{SELECT col FROM tab ORDER BY col LIMIT 1}. Index usage now happens automatically.

Move /contrib/pg_autovacuum into the main server (Alvaro)

Integrating autovacuum into the server allows it to be automatically started and stopped in sync with the database server, and allows autovacuum to be configured from \texttt{postgresql.conf}.

Add shared row level locks using \texttt{SELECT ... FOR SHARE} (Alvaro)

While PostgreSQL's MVCC locking allows \texttt{SELECT} to never be blocked by writers and therefore does not need shared row locks for typical operations, shared locks are useful for applications that require shared row locking. In particular this reduces the locking requirements imposed by referential integrity checks.

Add dependencies on shared objects, specifically roles (Alvaro)

This extension of the dependency mechanism prevents roles from being dropped while there are still database objects they own. Formerly it was possible to accidentally "orphan" objects by deleting their owner. While this could be recovered from, it was messy and unpleasant.

Improve performance for partitioned tables (Simon)

The new \texttt{constraint_exclusion} configuration parameter avoids lookups on child tables where constraints indicate that no matching rows exist in the child table.

This allows for a basic type of table partitioning. If child tables store separate key ranges and this is enforced using appropriate \texttt{CHECK} constraints, the optimizer will skip child table accesses when the constraint guarantees no matching rows exist in the child table.

**E.253.2. Migration to Version 8.1**

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

The 8.0 release announced that the \texttt{to_char()} function for intervals would be removed in 8.1. However, since no better API has been suggested, \texttt{to_char(interval)} has been enhanced in 8.1 and will remain in the server.

Observe the following incompatibilities:

- \texttt{add_missing_from} is now false by default (Neil)

  By default, we now generate an error if a table is used in a query without a FROM reference. The old behavior is still available, but the parameter must be set to 'true' to obtain it.

  It might be necessary to set \texttt{add_missing_from} to true in order to load an existing dump file, if the dump contains any views or rules created using the implicit-FROM syntax. This should be a one-time annoyance, because PostgreSQL 8.1 will convert such views and rules to standard explicit-FROM syntax. Subsequent dumps will therefore not have the problem.

- Cause input of a zero-length string ('') for float4/float8/oid to throw an error, rather than treating it as a zero (Neil)
This change is consistent with the current handling of zero-length strings for integers. The schedule for this change was announced in 8.0.

- **default_with_oids** is now false by default (Neil)

  With this option set to false, user-created tables no longer have an OID column unless **WITH OIDS** is specified in **CREATE TABLE**. Though OIDs have existed in all releases of PostgreSQL, their use is limited because they are only four bytes long and the counter is shared across all installed databases. The preferred way of uniquely identifying rows is via sequences and the **SERIAL** type, which have been supported since PostgreSQL 6.4.

- Add E'' syntax so eventually ordinary strings can treat backslashes literally (Bruce)

  Currently PostgreSQL processes a backslash in a string literal as introducing a special escape sequence, e.g. \n or \010. While this allows easy entry of special values, it is nonstandard and makes porting of applications from other databases more difficult. For this reason, the PostgreSQL project is planning to remove the special meaning of backslashes in strings. For backward compatibility and for users who want special backslash processing, a new string syntax has been created. This new string syntax is formed by writing an E immediately preceding the single quote that starts the string, e.g. E'hi\n'. While this release does not change the handling of backslashes in strings, it does add new configuration parameters to help users migrate applications for future releases:

  - **standard_conforming_strings** — does this release treat backslashes literally in ordinary strings?
  - **escape_string_warning** — warn about backslashes in ordinary (non-E) strings

The **standard_conforming_strings** value is read-only. Applications can retrieve the value to know how backslashes are processed. (Presence of the parameter can also be taken as an indication that E'' string syntax is supported.) In a future release, **standard_conforming_strings** will be true, meaning backslashes will be treated literally in non-E strings. To prepare for this change, use E'' strings in places that need special backslash processing, and turn on **escape_string_warning** to find additional strings that need to be converted to use E''.

  Also, use two single-quotes (''') to embed a literal single-quote in a string, rather than the PostgreSQL-supported syntax of backslash single-quote (\'). The former is standards-conforming and does not require the use of the E'' string syntax. You can also use the $$ string syntax, which does not treat backslashes specially.

- Make **REINDEX DATABASE** reindex all indexes in the database (Tom)

  Formerly, **REINDEX DATABASE** reindexed only system tables. This new behavior seems more intuitive. A new command **REINDEX SYSTEM** provides the old functionality of reindexing just the system tables.

- Read-only large object descriptors now obey MVCC snapshot semantics

  When a large object is opened with **INV_READ** (and not **INV_WRITE**), the data read from the descriptor will now reflect a “snapshot” of the large object's state at the time of the transaction snapshot in use by the query that called **lo_open()**. To obtain the old behavior of always returning the latest committed data, include **INV_WRITE** in the mode flags for **lo_open()**.

- Add proper dependencies for arguments of sequence functions (Tom)

  In previous releases, sequence names passed to **nextval()**, **currval()**, and **setval()** were stored as simple text strings, meaning that renaming or dropping a sequence used in a **DEFAULT** clause made the clause invalid. This release stores all newly-created sequence function arguments as internal OIDs, allowing them to track sequence renaming, and adding dependency information that prevents improper sequence removal. It also makes such **DEFAULT** clauses immune to schema renaming and search path changes.
Some applications might rely on the old behavior of run-time lookup for sequence names. This can still be done by explicitly casting the argument to text, for example nextval('myseq'::text).

Pre-8.1 database dumps loaded into 8.1 will use the old text-based representation and therefore will not have the features of OID-stored arguments. However, it is possible to update a database containing text-based DEFAULT clauses. First, save this query into a file, such as fixseq.sql:

```
SELECT 'ALTER TABLE ' ||
    pg_catalog.quote_ident(n.nspname) || '. ' ||
    pg_catalog.quote_ident(c.relname) ||
    ' ALTER COLUMN ' ||
    pg_catalog.quote_ident(a.attname) ||
    ' SET DEFAULT ' ||
    regexp_replace(d.adsrc,
        $$val\((\^[\']*')::text\)::regclass$$,
        $$val\(\1$$,
            'g') ||
    ';' FROM pg_namespace n, pg_class c, pg_attribute a, pg_attrdef d
WHERE n.oid = c.relnamespace AND c.oid = a.attrelid AND a.attrelid = d.adrelid AND a.attnum = d.adnum AND d.adsrc ~ $$val\((\^[\']*')::text\)::regclass$$;
```

Next, run the query against a database to find what adjustments are required, like this for database db1:

```
psql -t -f fixseq.sql db1
```

This will show the ALTER TABLE commands needed to convert the database to the newer OID-based representation. If the commands look reasonable, run this to update the database:

```
psql -t -f fixseq.sql db1 | psql -e db1
```

This process must be repeated in each database to be updated.

- In psql, treat unquoted \{digit\}+ sequences as octal (Bruce)

  In previous releases, \{digit\}+ sequences were treated as decimal, and only \0{digit}+ were treated as octal. This change was made for consistency.

- Remove grammar productions for prefix and postfix % and ^ operators (Tom)

  These have never been documented and complicated the use of the modulus operator (%) with negative numbers.

- Make &< and &> for polygons consistent with the box "over" operators (Tom)

- CREATE LANGUAGE can ignore the provided arguments in favor of information from pg_pltemplate (Tom)

  A new system catalog pg_pltemplate has been defined to carry information about the preferred definitions of procedural languages (such as whether they have validator functions). When an entry exists in this catalog for the language being created, CREATE LANGUAGE will ignore all its parameters except the language name and instead use the catalog information. This measure was taken because of increasing problems with obsolete language definitions being loaded by old dump
files. As of 8.1, pg_dump will dump procedural language definitions as just `CREATE LANGUAGE name`, relying on a template entry to exist at load time. We expect this will be a more future-proof representation.

- Make `pg_cancel_backend(int)` return a `boolean` rather than an `integer` (Neil)

- Some users are having problems loading UTF-8 data into 8.1.X. This is because previous versions allowed invalid UTF-8 byte sequences to be entered into the database, and this release properly accepts only valid UTF-8 sequences. One way to correct a dumpfile is to run the command `iconv -c -f UTF-8 -t UTF-8 -o cleanfile.sql dumpfile.sql`. The `-c` option removes invalid character sequences. A diff of the two files will show the sequences that are invalid. `iconv` reads the entire input file into memory so it might be necessary to use `split` to break up the dump into multiple smaller files for processing.

### E.253.3. Additional Changes

Below you will find a detailed account of the additional changes between PostgreSQL 8.1 and the previous major release.

#### E.253.3.1. Performance Improvements

- Improve GiST and R-tree index performance (Neil)

- Improve the optimizer, including auto-resizing of hash joins (Tom)

- Overhaul internal API in several areas

- Change WAL record CRCs from 64-bit to 32-bit (Tom)

  We determined that the extra cost of computing 64-bit CRCs was significant, and the gain in reliability too marginal to justify it.

- Prevent writing large empty gaps in WAL pages (Tom)

- Improve spinlock behavior on SMP machines, particularly Opterons (Tom)

- Allow nonconsecutive index columns to be used in a multicolumn index (Tom)

  For example, this allows an index on columns a,b,c to be used in a query with `WHERE a = 4 and c = 10`.

- Skip WAL logging for `CREATE TABLE AS / SELECT INTO` (Simon)

  Since a crash during `CREATE TABLE AS` would cause the table to be dropped during recovery, there is no reason to WAL log as the table is loaded. (Logging still happens if WAL archiving is enabled, however.)

- Allow concurrent GiST index access (Teodor, Oleg)

- Add configuration parameter `full_page_writes` to control writing full pages to WAL (Bruce)

  To prevent partial disk writes from corrupting the database, PostgreSQL writes a complete copy of each database disk page to WAL the first time it is modified after a checkpoint. This option turns off that functionality for more speed. This is safe to use with battery-backed disk caches where partial page writes cannot happen.

- Use `O_DIRECT` if available when using `O_SYNC` for `wal_sync_method` (Itagaki Takahiro)

  `O_DIRECT` causes disk writes to bypass the kernel cache, and for WAL writes, this improves performance.
• Improve COPY FROM performance (Alon Goldshuv)

   This was accomplished by reading COPY input in larger chunks, rather than character by character.

• Improve the performance of COUNT(), SUM, AVG(), STDDEV(), and VARIANCE () (Neil, Tom)

### E.253.3.2. Server Changes

• Prevent problems due to transaction ID (XID) wraparound (Tom)

   The server will now warn when the transaction counter approaches the wraparound point. If the counter becomes too close to wraparound, the server will stop accepting queries. This ensures that data is not lost before needed vacuuming is performed.

• Fix problems with object IDs (OIDs) conflicting with existing system objects after the OID counter has wrapped around (Tom)

• Add warning about the need to increase max_fsm_relations and max_fsm_pages during VACUUM (Ron Mayer)

• Add temp_buffers configuration parameter to allow users to determine the size of the local buffer area for temporary table access (Tom)

• Add session start time and client IP address to pg_stat_activity (Magnus)

• Adjust pg_stat views for bitmap scans (Tom)

   The meanings of some of the fields have changed slightly.

• Enhance pg_locks view (Tom)

• Log queries for client-side PREPARE and EXECUTE (Simon)

• Allow Kerberos name and user name case sensitivity to be specified in postgresql.conf (Magnus)

• Add configuration parameter krb_server_hostname so that the server host name can be specified as part of service principal (Todd Kover)

   If not set, any service principal matching an entry in the keytab can be used. This is new Kerberos matching behavior in this release.

• Add log_line_prefix options for millisecond timestamps (%m) and remote host (%h) (Ed L.)

• Add WAL logging for GiST indexes (Teodor, Oleg)

   GiST indexes are now safe for crash and point-in-time recovery.

• Remove old *.backup files when we do pg_stop_backup() (Bruce)

   This prevents a large number of *.backup files from existing in pg_xlog/.

• Add configuration parameters to control TCP/IP keep-alive times for idle, interval, and count (Oliver Jowett)

   These values can be changed to allow more rapid detection of lost client connections.

• Add per-user and per-database connection limits (Petr Jelinek)

   Using ALTER USER and ALTER DATABASE, limits can now be enforced on the maximum number of sessions that can concurrently connect as a specific user or to a specific database. Setting the limit to zero disables user or database connections.
• Allow more than two gigabytes of shared memory and per-backend work memory on 64-bit machines (Koichi Suzuki)

• New system catalog pg_pltemplate allows overriding obsolete procedural-language definitions in dump files (Tom)

E.253.3.3. Query Changes

• Add temporary views (Koju Iijima, Neil)

• Fix HAVING without any aggregate functions or GROUP BY so that the query returns a single group (Tom)

  Previously, such a case would treat the HAVING clause the same as a WHERE clause. This was not per spec.

• Add USING clause to allow additional tables to be specified to DELETE (Euler Taveira de Oliveira, Neil)

  In prior releases, there was no clear method for specifying additional tables to be used for joins in a DELETE statement. UPDATE already has a FROM clause for this purpose.

• Add support for \x hex escapes in backend and ecpg strings (Bruce)

  This is just like the standard C \x escape syntax. Octal escapes were already supported.

• Add BETWEEN SYMMETRIC query syntax (Pavel Stehule)

  This feature allows BETWEEN comparisons without requiring the first value to be less than the second. For example, 2 BETWEEN [ASYMMETRIC] 3 AND 1 returns false, while 2 BETWEEN SYMMETRIC 3 AND 1 returns true. BETWEEN ASYMMETRIC was already supported.

• Add NOWAIT option to SELECT ... FOR UPDATE/SHARE (Hans-Juergen Schoenig)

  While the statement_timeout configuration parameter allows a query taking more than a certain amount of time to be canceled, the NOWAIT option allows a query to be canceled as soon as a SELECT ... FOR UPDATE/SHARE command cannot immediately acquire a row lock.

E.253.3.4. Object Manipulation Changes

• Track dependencies of shared objects (Alvaro)

  PostgreSQL allows global tables (users, databases, tablespaces) to reference information in multiple databases. This addition adds dependency information for global tables, so, for example, user ownership can be tracked across databases, so a user who owns something in any database can no longer be removed. Dependency tracking already existed for database-local objects.

• Allow limited ALTER OWNER commands to be performed by the object owner (Stephen Frost)

  Prior releases allowed only superusers to change object owners. Now, ownership can be transferred if the user executing the command owns the object and would be able to create it as the new owner (that is, the user is a member of the new owning role and that role has the CREATE permission that would be needed to create the object afresh).

• Add ALTER object SET SCHEMA capability for some object types (tables, functions, types) (Bernd Helmle)

  This allows objects to be moved to different schemas.

• Add ALTER TABLE ENABLE/DISABLE TRIGGER to disable triggers (Satoshi Nagayasu)
E.253.3.5. Utility Command Changes

- Allow TRUNCATE to truncate multiple tables in a single command (Alvaro)

Because of referential integrity checks, it is not allowed to truncate a table that is part of a referential integrity constraint. Using this new functionality, TRUNCATE can be used to truncate such tables, if both tables involved in a referential integrity constraint are truncated in a single TRUNCATE command.

- Properly process carriage returns and line feeds in COPY CSV mode (Andrew)

In release 8.0, carriage returns and line feeds in CSV COPY TO were processed in an inconsistent manner. (This was documented on the TODO list.)

- Add COPY WITH CSV HEADER to allow a header line as the first line in COPY (Andrew)

This allows handling of the common CSV usage of placing the column names on the first line of the data file. For COPY TO, the first line contains the column names, and for COPY FROM, the first line is ignored.

- On Windows, display better sub-second precision in EXPLAIN ANALYZE (Magnus)

- Add trigger duration display to EXPLAIN ANALYZE (Tom)

Prior releases included trigger execution time as part of the total execution time, but did not show it separately. It is now possible to see how much time is spent in each trigger.

- Add support for \x hex escapes in COPY (Sergey Ten)

Previous releases only supported octal escapes.

- Make SHOW ALL include variable descriptions (Matthias Schmidt)

SHOW varname still only displays the variable's value and does not include the description.

- Make initdb create a new standard database called postgres, and convert utilities to use postgres rather than template1 for standard lookups (Dave)

In prior releases, template1 was used both as a default connection for utilities like createuser, and as a template for new databases. This caused CREATE DATABASE to sometimes fail, because a new database cannot be created if anyone else is in the template database. With this change, the default connection database is now postgres, meaning it is much less likely someone will be using template1 during CREATE DATABASE.

- Create new reindexdb command-line utility by moving /contrib/reindexdb into the server (Euler Taveira de Oliveira)

E.253.3.6. Data Type and Function Changes

- Add MAX() and MIN() aggregates for array types (Koju Iijima)

- Fix to_date() and to_timestamp() to behave reasonably when CC and YY fields are both used (Karel Zak)

  If the format specification contains CC and a year specification is YYY or longer, ignore the CC. If the year specification is YY or shorter, interpret CC as the previous century.

- Add md5(bytea) (Abhijit Menon-Sen)

  md5(text) already existed.

- Add support for numeric ^ numeric based on power(numeric, numeric)
The function already existed, but there was no operator assigned to it.

- **Fix NUMERIC modulus by properly truncating the quotient during computation (Bruce)**

  In previous releases, modulus for large values sometimes returned negative results due to rounding of the quotient.

- **Add a function lastval() (Dennis Björklund)**

  `lastval()` is a simplified version of `currval()`. It automatically determines the proper sequence name based on the most recent `nextval()` or `setval()` call performed by the current session.

- **Add `to_timestamp(DOUBLE PRECISION)` (Michael Glaesemann)**

  Converts Unix seconds since 1970 to a `TIMESTAMP WITH TIMEZONE`.

- **Add `pg_postmaster_start_time()` function (Euler Taveira de Oliveira, Matthias Schmidt)**

- **Allow the full use of time zone names in AT TIME ZONE, not just the short list previously available (Magnus)**

  Previously, only a predefined list of time zone names were supported by `AT TIME ZONE`. Now any supported time zone name can be used, e.g.:

  ```sql
  SELECT CURRENT_TIMESTAMP AT TIME ZONE 'Europe/London';
  ```

  In the above query, the time zone used is adjusted based on the daylight saving time rules that were in effect on the supplied date.

- **Add `GREATEST()` and `LEAST()` variadic functions (Pavel Stehule)**

  These functions take a variable number of arguments and return the greatest or least value among the arguments.

- **Add `pg_column_size()` (Mark Kirkwood)**

  This returns storage size of a column, which might be compressed.

- **Add `regexp_replace()` (Atsushi Ogawa)**

  This allows regular expression replacement, like sed. An optional flag argument allows selection of global (replace all) and case-insensitive modes.

- **Fix interval division and multiplication (Bruce)**

  Previous versions sometimes returned unjustified results, like `'4 months'::interval / 5` returning `'1 mon -6 days'`.

- **Fix roundoff behavior in timestamp, time, and interval output (Tom)**

  This fixes some cases in which the seconds field would be shown as 60 instead of incrementing the higher-order fields.

- **Add a separate day field to type `interval` so a one day interval can be distinguished from a 24 hour interval (Michael Glaesemann)**

  Days that contain a daylight saving time adjustment are not 24 hours long, but typically 23 or 25 hours. This change creates a conceptual distinction between intervals of “so many days” and intervals of “so many hours”. Adding `1 day` to a timestamp now gives the same local time on the
next day even if a daylight saving time adjustment occurs between, whereas adding 24 hours will give a different local time when this happens. For example, under US DST rules:

'2005-04-03 00:00:00-05' + '1 day' = '2005-04-04 00:00:00-04'
'2005-04-03 00:00:00-05' + '24 hours' = '2005-04-04 01:00:00-04'

• Add justify_days() and justify_hours() (Michael Glaesemann)

These functions, respectively, adjust days to an appropriate number of full months and days, and adjust hours to an appropriate number of full days and hours.

• Move /contrib/dbsize into the backend, and rename some of the functions (Dave Page, Andreas Pflug)
  • pg_tablespace_size()
  • pg_database_size()
  • pg_relation_size()
  • pg_total_relation_size()
  • pg_size_pretty()

pg_total_relation_size() includes indexes and TOAST tables.

• Add functions for read-only file access to the cluster directory (Dave Page, Andreas Pflug)
  • pg_stat_file()
  • pg_read_file()
  • pg_ls_dir()

• Add pg_reload_conf() to force reloading of the configuration files (Dave Page, Andreas Pflug)

• Add pg_rotate_logfile() to force rotation of the server log file (Dave Page, Andreas Pflug)

• Change pg_stat_* views to include TOAST tables (Tom)

E.253.3.7. Encoding and Locale Changes

• Rename some encodings to be more consistent and to follow international standards (Bruce)
  • UNICODE is now UTF8
  • ALT is now WIN866
  • WIN is now WIN1251
  • TCVN is now WIN1258

The original names still work.

• Add support for WIN1252 encoding (Roland Volkmann)

• Add support for four-byte UTF8 characters (John Hansen)

Previously only one, two, and three-byte UTF8 characters were supported. This is particularly important for support for some Chinese character sets.
• Allow direct conversion between EUC_JP and SJIS to improve performance (Atushi Ogawa)

• Allow the UTF8 encoding to work on Windows (Magnus)

This is done by mapping UTF8 to the Windows-native UTF16 implementation.

**E.253.3.8. General Server-Side Language Changes**

• Fix ALTER LANGUAGE RENAME (Sergey Yatskevich)

• Allow function characteristics, like strictness and volatility, to be modified via ALTER FUNCTION (Neil)

• Increase the maximum number of function arguments to 100 (Tom)

• Allow SQL and PL/pgSQL functions to use OUT and INOUT parameters (Tom)

OUT is an alternate way for a function to return values. Instead of using RETURN, values can be returned by assigning to parameters declared as OUT or INOUT. This is notationally simpler in some cases, particularly so when multiple values need to be returned. While returning multiple values from a function was possible in previous releases, this greatly simplifies the process. (The feature will be extended to other server-side languages in future releases.)

• Move language handler functions into the pg_catalog schema

This makes it easier to drop the public schema if desired.

• Add SPI_getnspname() to SPI (Neil)

**E.253.3.9. PL/pgSQL Server-Side Language Changes**

• Overhaul the memory management of PL/pgSQL functions (Neil)

The parsetree of each function is now stored in a separate memory context. This allows this memory to be easily reclaimed when it is no longer needed.

• Check function syntax at CREATE FUNCTION time, rather than at runtime (Neil)

Previously, most syntax errors were reported only when the function was executed.

• Allow OPEN to open non-SELECT queries like EXPLAIN and SHOW (Tom)

• No longer require functions to issue a RETURN statement (Tom)

This is a byproduct of the newly added OUT and INOUT functionality. RETURN can be omitted when it is not needed to provide the function's return value.

• Add support for an optional INTO clause to PL/pgSQL's EXECUTE statement (Pavel Stehule, Neil)

• Make CREATE TABLE AS set ROW_COUNT (Tom)

• Define SQLSTATE and SQLERRM to return the SQLSTATE and error message of the current exception (Pavel Stehule, Neil)

These variables are only defined inside exception blocks.

• Allow the parameters to the RAISE statement to be expressions (Pavel Stehule, Neil)

• Add a loop CONTINUE statement (Pavel Stehule, Neil)

• Allow block and loop labels (Pavel Stehule)
E.253.3.10. PL/Perl Server-Side Language Changes

- Allow large result sets to be returned efficiently (Abhijit Menon-Sen)
  This allows functions to use `return_next()` to avoid building the entire result set in memory.
- Allow one-row-at-a-time retrieval of query results (Abhijit Menon-Sen)
  This allows functions to use `spi_query()` and `spi_fetchrow()` to avoid accumulating the entire result set in memory.
- Force PL/Perl to handle strings as UTF8 if the server encoding is UTF8 (David Kamholz)
- Add a validator function for PL/Perl (Andrew)
  This allows syntax errors to be reported at definition time, rather than execution time.
- Allow PL/Perl to return a Perl array when the function returns an array type (Andrew)
  This basically maps PostgreSQL arrays to Perl arrays.
- Allow Perl nonfatal warnings to generate NOTICE messages (Andrew)
- Allow Perl's `strict` mode to be enabled (Andrew)

E.253.3.11. psql Changes

- Add `\set ON_ERROR_ROLLBACK` to allow statements in a transaction to error without affecting the rest of the transaction (Greg Sabino Mullane)
  This is basically implemented by wrapping every statement in a sub-transaction.
- Add support for `\x` hex strings in psql variables (Bruce)
  Octal escapes were already supported.
- Add support for `troff -ms` output format (Roger Leigh)
- Allow the history file location to be controlled by `HISTFILE` (Andreas Seltenreich)
  This allows configuration of per-database history storage.
- Prevent `\x` (expanded mode) from affecting the output of `\d tablename` (Neil)
- Add `-L` option to psql to log sessions (Lorne Sunley)
  This option was added because some operating systems do not have simple command-line activity logging functionality.
- Make `\d` show the tablespaces of indexes (Qingqing Zhou)
- Allow psql help (`\h`) to make a best guess on the proper help information (Greg Sabino Mullane)
  This allows the user to just add `\h` to the front of the syntax error query and get help on the supported syntax. Previously any additional query text beyond the command name had to be removed to use `\h`.
- Add `\pset numericlocale` to allow numbers to be output in a locale-aware format (Eugen Nedelcu)
  For example, using `C` locale `100000` would be output as `100,000.0` while a European locale might output this value as `100.000,0`. 

3117
• Make startup banner show both server version number and psql's version number, when they are different (Bruce)

Also, a warning will be shown if the server and psql are from different major releases.

E.253.3.12. pg_dump Changes

• Add \-n /\--schema switch to pg_restore (Richard van den Berg)

This allows just the objects in a specified schema to be restored.

• Allow pg_dump to dump large objects even in text mode (Tom)

With this change, large objects are now always dumped; the former \-b switch is a no-op.

• Allow pg_dump to dump a consistent snapshot of large objects (Tom)

• Dump comments for large objects (Tom)

• Add \--encoding to pg_dump (Magnus Hagander)

This allows a database to be dumped in an encoding that is different from the server's encoding. This is valuable when transferring the dump to a machine with a different encoding.

• Rely on pg_pltemplate for procedural languages (Tom)

If the call handler for a procedural language is in the pg_catalog schema, pg_dump does not dump the handler. Instead, it dumps the language using just CREATE LANGUAGE name, relying on the pg_pltemplate catalog to provide the language's creation parameters at load time.

E.253.3.13. libpq Changes

• Add a PGPASSFILE environment variable to specify the password file's filename (Andrew)

• Add lo_create(), that is similar to lo_creat() but allows the OID of the large object to be specified (Tom)

• Make libpq consistently return an error to the client application on malloc() failure (Neil)

E.253.3.14. Source Code Changes

• Fix pgxs to support building against a relocated installation

• Add spinlock support for the Itanium processor using Intel compiler (Vikram Kalsi)

• Add Kerberos 5 support for Windows (Magnus)

• Add Chinese FAQ (laser@pgsqldb.com)

• Rename Rendezvous to Bonjour to match OS/X feature renaming (Bruce)

• Add support for fsync_writethrough on macOS (Chris Campbell)

• Streamline the passing of information within the server, the optimizer, and the lock system (Tom)

• Allow pg_config to be compiled using MSVC (Andrew)

This is required to build DBD::Pg using MSVC.

• Remove support for Kerberos V4 (Magnus)
Kerberos 4 had security vulnerabilities and is no longer maintained.

- Code cleanups (Coverity static analysis performed by EnterpriseDB)
- Modify postgresql.conf to use documentation defaults on/off rather than true/false (Bruce)
- Enhance pg_config to be able to report more build-time values (Tom)
- Allow libpq to be built thread-safe on Windows (Dave Page)
- Allow IPv6 connections to be used on Windows (Andrew)
- Add Server Administration documentation about I/O subsystem reliability (Bruce)
- Move private declarations from gist.h to gist_private.h (Neil)

In previous releases, gist.h contained both the public GiST API (intended for use by authors of GiST index implementations) as well as some private declarations used by the implementation of GiST itself. The latter have been moved to a separate file, gist_private.h. Most GiST index implementations should be unaffected.

- Overhaul GiST memory management (Neil)

GiST methods are now always invoked in a short-lived memory context. Therefore, memory allocated via palloc() will be reclaimed automatically, so GiST index implementations do not need to manually release allocated memory via pfree().

E.253.3.15. Contrib Changes

- Add /contrib/pg_buffercache contrib module (Mark Kirkwood)
  This displays the contents of the buffer cache, for debugging and performance tuning purposes.
- Remove /contrib/array because it is obsolete (Tom)
- Clean up the /contrib/lo module (Tom)
- Move /contrib/findoidjoins to /src/tools (Tom)
- Remove the <<,>>, &<, and &> operators from /contrib/cube
  These operators were not useful.
- Improve /contrib/btree_gist (Janko Richter)
- Improve /contrib/pgbench (Tomoaki Sato, Tatsuo)
  There is now a facility for testing with SQL command scripts given by the user, instead of only a hard-wired command sequence.
- Improve /contrib/pgcrypto (Marko Kreen)
  Implementation of OpenPGP symmetric-key and public-key encryption
  Both RSA and Elgamal public-key algorithms are supported.
- Stand alone build: include SHA256/384/512 hashes, Fortuna PRNG
- OpenSSL build: support 3DES, use internal AES with OpenSSL < 0.9.7
• Take build parameters (OpenSSL, zlib) from configure result

There is no need to edit the Makefile anymore.

• Remove support for libmhash and libmcrypt

E.254. Release 8.0.26

Release date: 2010-10-04

This release contains a variety of fixes from 8.0.25. For information about new features in the 8.0 major release, see Section E.280.

This is expected to be the last PostgreSQL release in the 8.0.X series. Users are encouraged to update to a newer release branch soon.

E.254.1. Migration to Version 8.0.26

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.22, see Section E.258.

E.254.2. Changes

• Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function's owner.

The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

• Prevent possible crashes in pg_get_expr() by disallowing it from being called with an argument that is not one of the system catalog columns it's intended to be used with (Heikki Linnakangas, Tom Lane)

• Fix "cannot handle unplanned sub-select" error (Tom Lane)

This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.

• Defend against functions returning setof record where not all the returned rows are actually of the same rowtype (Tom Lane)

• Take care to fsync the contents of lockfiles (both postmaster.pid and the socket lockfile) while writing them (Tom Lane)
This omission could result in corrupted lockfile contents if the machine crashes shortly after postmaster start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.

- Avoid recursion while assigning XIDs to heavily-nested subtransactions (Andres Freund, Robert Haas)

  The original coding could result in a crash if there was limited stack space.

- Fix \texttt{log\_line\_prefix}'s \%i escape, which could produce junk early in backend startup (Tom Lane)

- Fix possible data corruption in \texttt{ALTER TABLE ... SET TABLESPACE} when archiving is enabled (Jeff Davis)

- Allow \texttt{CREATE DATABASE} and \texttt{ALTER DATABASE ... SET TABLESPACE} to be interrupted by query-cancel (Guillaume Lelarge)

- In PL/Python, defend against null pointer results from \texttt{PyCObject\_AsVoidPtr} and \texttt{PyCObject\_FromVoidPtr} (Peter Eisentraut)

- Improve contrib/dblink's handling of tables containing dropped columns (Tom Lane)

- Fix connection leak after "duplicate connection name" errors in contrib/dblink (Itagaki Takahiro)

- Fix contrib/dblink to handle connection names longer than 62 bytes correctly (Itagaki Takahiro)

- Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)

- Update time zone data files to tzdata release 2010l for DST law changes in Egypt and Palestine; also historical corrections for Finland.

  This change also adds new names for two Micronesian timezones: Pacific/Chuuk is now preferred over Pacific/Truk (and the preferred abbreviation is CHUT not TRUT) and Pacific/Pohnpei is preferred over Pacific/Ponape.

\textbf{E.255. Release 8.0.25}

\textbf{Release date:} 2010-05-17

This release contains a variety of fixes from 8.0.24. For information about new features in the 8.0 major release, see Section E.280.

The PostgreSQL community will stop releasing updates for the 8.0.X release series in July 2010. Users are encouraged to update to a newer release branch soon.

\textbf{E.255.1. Migration to Version 8.0.25}

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.22, see Section E.258.

\textbf{E.255.2. Changes}

- Enforce restrictions in \texttt{plperl} using an opmask applied to the whole interpreter, instead of using \texttt{Safe.pm} (Tim Bunce, Andrew Dunstan)
Recent developments have convinced us that Safe.pm is too insecure to rely on for making plperl trustable. This change removes use of Safe.pm altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include that it is now possible to use Perl's strict pragma in a natural way in plperl, and that Perl's $a and $b variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

- Prevent PL/Tcl from executing untrustworthy code from pltcl_modules (Tom)

  PL/Tcl's feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless pltcl_modules is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted "normal" Tcl interpreter unless we are really going to execute a pltclu function. (CVE-2010-1170)

- Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

  Previously, if an unprivileged user ran ALTER USER ... RESET ALL for himself, or ALTER DATABASE ... RESET ALL for a database he owns, this would remove all special parameter settings for the user or database, even ones that are only supposed to be changeable by a superuser. Now, the ALTER will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a CONTEXT addition would be made to log entries (Tom)

  In some cases the context-printing function would fail because the current transaction had already been rolled back when it came time to print a log message.

- Update PL/Perl's ppport.h for modern Perl versions (Andrew)

- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)

- Prevent infinite recursion in psql when expanding a variable that refers to itself (Tom)

- Ensure that contrib/pgstattuple functions respond to cancel interrupts promptly (Tatsuhito Kasahara)

- Make server startup deal properly with the case that shmget() returns EINVAL for an existing shared memory segment (Tom)

  This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entirely-misleading startup failure complaining that the shared memory request size was too large.

- Update time zone data files to tzdata release 2010j for DST law changes in Argentina, Australian Antarctic, Bangladesh, Mexico, Morocco, Pakistan, Palestine, Russia, Syria, Tunisia; also historical corrections for Taiwan.

E.256. Release 8.0.24

Release date: 2010-03-15

This release contains a variety of fixes from 8.0.23. For information about new features in the 8.0 major release, see Section E.280.

The PostgreSQL community will stop releasing updates for the 8.0.X release series in July 2010. Users are encouraged to update to a newer release branch soon.
E.256.1. Migration to Version 8.0.24

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.22, see Section E.258.

E.256.2. Changes

• Add new configuration parameter `ssl_renegotiation_limit` to control how often we do session key renegotiation for an SSL connection (Magnus)

  This can be set to zero to disable renegotiation completely, which may be required if a broken SSL library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that cause renegotiation attempts to fail.

• Fix possible crashes when trying to recover from a failure in subtransaction start (Tom)

• Fix server memory leak associated with use of savepoints and a client encoding different from server's encoding (Tom)

• Make `substring()` for bit types treat any negative length as meaning “all the rest of the string” (Tom)

  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

• Fix integer-to-bit-string conversions to handle the first fractional byte correctly when the output bit width is wider than the given integer by something other than a multiple of 8 bits (Tom)

• Fix some cases of pathologically slow regular expression matching (Tom)

• Fix the `STOP WAL LOCATION` entry in backup history files to report the next WAL segment's name when the end location is exactly at a segment boundary (Itagaki Takahiro)

• When reading `pg_hba.conf` and related files, do not treat `@something` as a file inclusion request if the `@` appears inside quote marks; also, never treat `@` by itself as a file inclusion request (Tom)

  This prevents erratic behavior if a role or database name starts with `@`. If you need to include a file whose path name contains spaces, you can still do so, but you must write `@"/path to/file"` rather than putting the quotes around the whole construct.

• Prevent infinite loop on some platforms if a directory is named as an inclusion target in `pg_hba.conf` and related files (Tom)

• Fix pgsql failure in one case where a composite column is set to NULL (Tom)

• Add `volatile` markings in PL/Python to avoid possible compiler-specific misbehavior (Zdenek Kotala)

• Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)

  The only known symptom of this oversight is that the Tcl `clock` command misbehaves if using Tcl 8.5 or later.

• Prevent crash in `contrib/dblink` when too many key columns are specified to a `dblink_build_sql_*` function (Rushabh Lathia, Joe Conway)

• Fix assorted crashes in `contrib/xml2` caused by sloppy memory management (Tom)

• Update time zone data files to tzdata release 2010e for DST law changes in Bangladesh, Chile, Fiji, Mexico, Paraguay, Samoa.
E.257. Release 8.0.23

Release date: 2009-12-14

This release contains a variety of fixes from 8.0.22. For information about new features in the 8.0 major release, see Section E.280.

E.257.1. Migration to Version 8.0.23

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.22, see Section E.258.

E.257.2. Changes

- Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)
  
  This change prevents allegedly-immutable index functions from possibly subverting a superuser's session (CVE-2009-4136).

- Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)
  
  This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).

- Fix possible crash during backend-startup-time cache initialization (Tom)

- Prevent signals from interrupting VACUUM at unsafe times (Alvaro)
  
  This fix prevents a PANIC if a VACUUM FULL is canceled after it's already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

- Fix possible crash due to integer overflow in hash table size calculation (Tom)
  
  This could occur with extremely large planner estimates for the size of a hashjoin's result.

- Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)

- Fix premature drop of temporary files used for a cursor that is accessed within a subtransaction (Heikki)

- Fix PAM password processing to be more robust (Tom)
  
  The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.

- Fix rare crash in exception processing in PL/Python (Peter)

- Ensure psql's flex module is compiled with the correct system header definitions (Tom)
  
  This fixes build failures on platforms where --enable-largefile causes incompatible changes in the generated code.

- Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

- Update time zone data files to tzdata release 2009s for DST law changes in Antarctica, Argentina, Bangladesh, Fiji, Novokuznetsk, Pakistan, Palestine, Samoa, Syria; also historical corrections for Hong Kong.
E.258. Release 8.0.22

Release date: 2009-09-09

This release contains a variety of fixes from 8.0.21. For information about new features in the 8.0 major release, see Section E.280.

E.258.1. Migration to Version 8.0.22

A dump/restore is not required for those running 8.0.X. However, if you have any hash indexes on interval columns, you must REINDEX them after updating to 8.0.22. Also, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.258.2. Changes

- Disallow RESET ROLE and RESET SESSION AUTHORIZATION inside security-definer functions (Tom, Heikki)

  This covers a case that was missed in the previous patch that disallowed SET ROLE and SET SESSION AUTHORIZATION inside security-definer functions. (See CVE-2007-6600)

- Fix handling of sub-SELECTs appearing in the arguments of an outer-level aggregate function (Tom)

- Fix hash calculation for data type interval (Tom)

  This corrects wrong results for hash joins on interval values. It also changes the contents of hash indexes on interval columns. If you have any such indexes, you must REINDEX them after updating.

- Treat to_char(..., 'TH') as an uppercase ordinal suffix with 'HH'/HH12' (Heikki)

  It was previously handled as 'th' (lowercase).

- Fix overflow for INTERVAL 'x ms' when x is more than 2 million and integer datetimes are in use (Alex Hunsaker)

- Fix calculation of distance between a point and a line segment (Tom)

  This led to incorrect results from a number of geometric operators.

- Fix money data type to work in locales where currency amounts have no fractional digits, e.g. Japan (Itagaki Takahiro)

- Properly round datetime input like 00:12:57.999999999999999999999999999999 (Tom)

- Fix poor choice of page split point in GiST R-tree operator classes (Teodor)

- Fix portability issues in plperl initialization (Andrew Dunstan)

- Fix pg_ctl to not go into an infinite loop if postgresql.conf is empty (Jeff Davis)

- Fix contrib/xml2's xslt_process() to properly handle the maximum number of parameters (twenty) (Tom)

- Improve robustness of libpq's code to recover from errors during COPY FROM STDIN (Tom)

- Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)
• Update time zone data files to tzdata release 2009l for DST law changes in Bangladesh, Egypt, Jordan, Pakistan, Argentina/San_Luis, Cuba, Jordan (historical correction only), Mauritius, Morocco, Palestine, Syria, Tunisia.

E.259. Release 8.0.21

Release date: 2009-03-16

This release contains a variety of fixes from 8.0.20. For information about new features in the 8.0 major release, see Section E.280.

E.259.1. Migration to Version 8.0.21

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.259.2. Changes

• Prevent error recursion crashes when encoding conversion fails (Tom)

This change extends fixes made in the last two minor releases for related failure scenarios. The previous fixes were narrowly tailored for the original problem reports, but we have now recognized that any error thrown by an encoding conversion function could potentially lead to infinite recursion while trying to report the error. The solution therefore is to disable translation and encoding conversion and report the plain-ASCII form of any error message, if we find we have gotten into a recursive error reporting situation. (CVE-2009-0922)

• Disallow CREATE CONVERSION with the wrong encodings for the specified conversion function (Heikki)

This prevents one possible scenario for encoding conversion failure. The previous change is a backstop to guard against other kinds of failures in the same area.

• Fix core dump when to_char() is given format codes that are inappropriate for the type of the data argument (Tom)

• Add MUST (Mauritius Island Summer Time) to the default list of known timezone abbreviations (Xavier Bugaud)

E.260. Release 8.0.20

Release date: 2009-02-02

This release contains a variety of fixes from 8.0.19. For information about new features in the 8.0 major release, see Section E.280.

E.260.1. Migration to Version 8.0.20

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.260.2. Changes

• Improve handling of URLs in headline() function (Teodor)

• Improve handling of overlength headlines in headline() function (Teodor)
• Prevent possible Assert failure or misconversion if an encoding conversion is created with the wrong conversion function for the specified pair of encodings (Tom, Heikki)

• Avoid unnecessary locking of small tables in VACUUM (Heikki)

• Fix uninitialized variables in contrib/tsearch2's get_covers() function (Teodor)

• Make all documentation reference postgresql-bugs and/or postgresql-hackers as appropriate, instead of the now-decommissioned postgresql-ports and postgresql-patches mailing lists (Tom)

• Update time zone data files to tzdata release 2009a (for Kathmandu and historical DST corrections in Switzerland, Cuba)

E.261. Release 8.0.19

Release date: 2008-11-03

This release contains a variety of fixes from 8.0.18. For information about new features in the 8.0 major release, see Section E.280.

E.261.1. Migration to Version 8.0.19

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.261.2. Changes

• Fix backend crash when the client encoding cannot represent a localized error message (Tom)

  We have addressed similar issues before, but it would still fail if the “character has no equivalent” message itself couldn't be converted. The fix is to disable localization and send the plain ASCII error message when we detect such a situation.

• Fix possible crash when deeply nested functions are invoked from a trigger (Tom)

• Ensure an error is reported when a newly-defined PL/pgSQL trigger function is invoked as a normal function (Tom)

• Fix incorrect tsearch2 headline generation when single query item matches first word of text (Sushant Sinha)

• Fix improper display of fractional seconds in interval values when using a non-ISO datestyle in an --enable-integer-datetimes build (Ron Mayer)

• Ensure SPI_getvalue and SPI_getbinval behave correctly when the passed tuple and tuple descriptor have different numbers of columns (Tom)

  This situation is normal when a table has had columns added or removed, but these two functions didn't handle it properly. The only likely consequence is an incorrect error indication.

• Fix ecpg's parsing of CREATE USER (Michael)

• Fix recent breakage of pg_ctl restart (Tom)

• Update time zone data files to tzdata release 2008i (for DST law changes in Argentina, Brazil, Mauritius, Syria)

E.262. Release 8.0.18

Release date: 2008-09-22
This release contains a variety of fixes from 8.0.17. For information about new features in the 8.0 major release, see Section E.280.

E.262.1. Migration to Version 8.0.18

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.262.2. Changes

- Widen local lock counters from 32 to 64 bits (Tom)

  This responds to reports that the counters could overflow in sufficiently long transactions, leading to unexpected “lock is already held” errors.

- Add checks in executor startup to ensure that the tuples produced by an INSERT or UPDATE will match the target table’s current rowtype (Tom)

  ALTER COLUMN TYPE, followed by re-use of a previously cached plan, could produce this type of situation. The check protects against data corruption and/or crashes that could ensue.

- Fix datetime input functions to correctly detect integer overflow when running on a 64-bit platform (Tom)

- Improve performance of writing very long log messages to syslog (Tom)

- Fix bug in backwards scanning of a cursor on a SELECT DISTINCT ON query (Tom)

- Fix planner to estimate that GROUP BY expressions yielding boolean results always result in two groups, regardless of the expressions’ contents (Tom)

  This is very substantially more accurate than the regular GROUP BY estimate for certain boolean tests like col IS NULL.

- Fix PL/Tcl to behave correctly with Tcl 8.5, and to be more careful about the encoding of data sent to or from Tcl (Tom)

- Fix PL/Python to work with Python 2.5

  This is a back-port of fixes made during the 8.2 development cycle.

- Improve pg_dump and pg_restore’s error reporting after failure to send a SQL command (Tom)

- Fix pg_ctl to properly preserve postmaster command-line arguments across a restart (Bruce)

- Update time zone data files to tzdata release 2008f (for DST law changes in Argentina, Bahamas, Brazil, Mauritius, Morocco, Pakistan, Palestine, and Paraguay)

E.263. Release 8.0.17

Release date: 2008-06-12

This release contains one serious bug fix over 8.0.16. For information about new features in the 8.0 major release, see Section E.280.

E.263.1. Migration to Version 8.0.17

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.
E.263.2. Changes

• Make `pg_get_ruledef()` parenthesize negative constants (Tom)

Before this fix, a negative constant in a view or rule might be dumped as, say, `-42::integer`, which is subtly incorrect: it should be `(-42)::integer` due to operator precedence rules. Usually this would make little difference, but it could interact with another recent patch to cause PostgreSQL to reject what had been a valid `SELECT DISTINCT` view query. Since this could result in `pg_dump` output failing to reload, it is being treated as a high-priority fix. The only released versions in which dump output is actually incorrect are 8.3.1 and 8.2.7.

E.264. Release 8.0.16

**Release date:** never released

This release contains a variety of fixes from 8.0.15. For information about new features in the 8.0 major release, see Section E.280.

E.264.1. Migration to Version 8.0.16

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.264.2. Changes

• Fix `ALTER TABLE ADD COLUMN ... PRIMARY KEY` so that the new column is correctly checked to see if it's been initialized to all non-nulls (Brendan Jurd)

Previous versions neglected to check this requirement at all.

• Fix possible `CREATE TABLE` failure when inheriting the “same” constraint from multiple parent relations that inherited that constraint from a common ancestor (Tom)

• Fix conversions between ISO-8859-5 and other encodings to handle Cyrillic “Yo” characters (e and E with two dots) (Sergey Burladyan)

• Fix a few datatype input functions that were allowing unused bytes in their results to contain uninitialized, unpredictable values (Tom)

This could lead to failures in which two apparently identical literal values were not seen as equal, resulting in the parser complaining about unmatched `ORDER BY` and `DISTINCT` expressions.

• Fix a corner case in regular-expression substring matching (`substring(string from pattern)`) (Tom)

The problem occurs when there is a match to the pattern overall but the user has specified a parenthesized subexpression and that subexpression hasn't got a match. An example is `substring('foo' from 'foo(bar)?')`. This should return NULL, since `bar` isn't matched, but it was mistakenly returning the whole-pattern match instead (ie, `foo`).

• Update time zone data files to tzdata release 2008c (for DST law changes in Morocco, Iraq, Choibalsan, Pakistan, Syria, Cuba, Argentina/San_Luis, and Chile)

• Fix incorrect result from ecpg's `PGTYPEStimestamp_sub()` function (Michael)

• Fix core dump in contrib/xml2's `xpath_table()` function when the input query returns a NULL value (Tom)
• Fix contrib/xml2's makefile to not override CFLAGS (Tom)

• Fix DatumGetBool macro to not fail with gcc 4.3 (Tom)

This problem affects "old style" (V0) C functions that return boolean. The fix is already in 8.3, but the need to back-patch it was not realized at the time.

• Fix longstanding LISTEN/NOTIFY race condition (Tom)

In rare cases a session that had just executed a LISTEN might not get a notification, even though one would be expected because the concurrent transaction executing NOTIFY was observed to commit later.

A side effect of the fix is that a transaction that has executed a not-yet-committed LISTEN command will not see any row in pg_listener for the LISTEN, should it choose to look; formerly it would have. This behavior was never documented one way or the other, but it is possible that some applications depend on the old behavior.

• Fix rare crash when an error occurs during a query using a hash index (Heikki)

• Fix input of datetime values for February 29 in years BC (Tom)

The former coding was mistaken about which years were leap years.

• Fix "unrecognized node type" error in some variants of ALTER OWNER (Tom)

• Fix pg_ctl to correctly extract the postmaster's port number from command-line options (Itagaki Takahiro, Tom)

Previously, pg_ctl start -w could try to contact the postmaster on the wrong port, leading to bogus reports of startup failure.

• Use -fwrapv to defend against possible misoptimization in recent gcc versions (Tom)

This is known to be necessary when building PostgreSQL with gcc 4.3 or later.

• Fix display of constant expressions in ORDER BY and GROUP BY (Tom)

An explicitly casted constant would be shown incorrectly. This could for example lead to corruption of a view definition during dump and reload.

• Fix libpq to handle NOTICE messages correctly during COPY OUT (Tom)

This failure has only been observed to occur when a user-defined datatype's output routine issues a NOTICE, but there is no guarantee it couldn't happen due to other causes.

E.265. Release 8.0.15

Release date: 2008-01-07

This release contains a variety of fixes from 8.0.14, including fixes for significant security issues. For information about new features in the 8.0 major release, see Section E.280.

This is the last 8.0.X release for which the PostgreSQL community will produce binary packages for Windows. Windows users are encouraged to move to 8.2.X or later, since there are Windows-specific fixes in 8.2.X that are impractical to back-port. 8.0.X will continue to be supported on other platforms.

E.265.1. Migration to Version 8.0.15

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.
E.265.2. Changes

- Prevent functions in indexes from executing with the privileges of the user running VACUUM, ANALYZE, etc (Tom)

Functions used in index expressions and partial-index predicates are evaluated whenever a new table entry is made. It has long been understood that this poses a risk of trojan-horse code execution if one modifies a table owned by an untrustworthy user. (Note that triggers, defaults, check constraints, etc. pose the same type of risk.) But functions in indexes pose extra danger because they will be executed by routine maintenance operations such as VACUUM FULL, which are commonly performed automatically under a superuser account. For example, a nefarious user can execute code with superuser privileges by setting up a trojan-horse index definition and waiting for the next routine vacuum. The fix arranges for standard maintenance operations (including VACUUM, ANALYZE, REINDEX, and CLUSTER) to execute as the table owner rather than the calling user, using the same privilege-switching mechanism already used for SECURITY DEFINER functions. To prevent bypassing this security measure, execution of SET SESSION AUTHORIZATION and SET ROLE is now forbidden within a SECURITY DEFINER context. (CVE-2007-6600)

- Repair assorted bugs in the regular-expression package (Tom, Will Drewry)

Suitably crafted regular-expression patterns could cause crashes, infinite or near-infinite looping, and/or massive memory consumption, all of which pose denial-of-service hazards for applications that accept regex search patterns from untrustworthy sources. (CVE-2007-4769, CVE-2007-4772, CVE-2007-6067)

- Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

The fix that appeared for this in 8.0.14 was incomplete, as it plugged the hole for only some dblink functions. (CVE-2007-6601, CVE-2007-3278)

- Update time zone data files to tzdata release 2007k (in particular, recent Argentina changes) (Tom)

- Fix planner failure in some cases of WHERE false AND var IN (SELECT ...) (Tom)

- Preserve the tablespace of indexes that are rebuilt by ALTER TABLE ... ALTER COLUMN TYPE (Tom)

- Make archive recovery always start a new WAL timeline, rather than only when a recovery stop time was used (Simon)

This avoids a corner-case risk of trying to overwrite an existing archived copy of the last WAL segment, and seems simpler and cleaner than the original definition.

- Make VACUUM not use all of maintenance_work_mem when the table is too small for it to be useful (Alvaro)

- Fix potential crash in translate() when using a multibyte database encoding (Tom)

- Fix PL/Perl to cope when platform's Perl defines type bool as int rather than char (Tom)

While this could theoretically happen anywhere, no standard build of Perl did things this way ... until macOS 10.5.

- Fix PL/Python to not crash on long exception messages (Alvaro)

- Fix pg_dump to correctly handle inheritance child tables that have default expressions different from their parent's (Tom)

- ecpg parser fixes (Michael)
• Make contrib/tablefunc's crosstab() handle NULL rowid as a category in its own right, rather than crashing (Joe)

• Fix tsvector and tsquery output routines to escape backslashes correctly (Teodor, Bruce)

• Fix crash of to_tsvector() on huge input strings (Teodor)

• Require a specific version of Autoconf to be used when re-generating the configure script (Peter)

This affects developers and packagers only. The change was made to prevent accidental use of untested combinations of Autoconf and PostgreSQL versions. You can remove the version check if you really want to use a different Autoconf version, but it's your responsibility whether the result works or not.

E.266. Release 8.0.14

Release date:  2007-09-17

This release contains a variety of fixes from 8.0.13. For information about new features in the 8.0 major release, see Section E.280.

E.266.1. Migration to Version 8.0.14

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.266.2. Changes

• Prevent index corruption when a transaction inserts rows and then aborts close to the end of a concurrent VACUUM on the same table (Tom)

• Make CREATE DOMAIN ... DEFAULT NULL work properly (Tom)

• Fix excessive logging of SSL error messages (Tom)

• Fix logging so that log messages are never interleaved when using the syslogger process (Andrew)

• Fix crash when log_min_error_statement logging runs out of memory (Tom)

• Fix incorrect handling of some foreign-key corner cases (Tom)

• Prevent CLUSTER from failing due to attempting to process temporary tables of other sessions (Alvaro)

• Update the time zone database rules, particularly New Zealand's upcoming changes (Tom)

• Windows socket improvements (Magnus)

• Suppress timezone name (%Z) in log timestamps on Windows because of possible encoding mismatches (Tom)

• Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

E.267. Release 8.0.13

Release date:  2007-04-23

This release contains a variety of fixes from 8.0.12, including a security fix. For information about new features in the 8.0 major release, see Section E.280.
E.267.1. Migration to Version 8.0.13

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.267.2. Changes

• Support explicit placement of the temporary-table schema within search_path, and disable searching it for functions and operators (Tom)

  This is needed to allow a security-definer function to set a truly secure value of search_path. Without it, an unprivileged SQL user can use temporary objects to execute code with the privileges of the security-definer function (CVE-2007-2138). See CREATE FUNCTION for more information.

• /contrib/tsearch2 crash fixes (Teodor)

• Fix potential-data-corruption bug in how VACUUM FULL handles UPDATE chains (Tom, Pavan Deolasee)

• Fix PANIC during enlargement of a hash index (bug introduced in 8.0.10) (Tom)

• Fix POSIX-style timezone specs to follow new USA DST rules (Tom)

E.268. Release 8.0.12

Release date: 2007-02-07

This release contains one fix from 8.0.11. For information about new features in the 8.0 major release, see Section E.280.

E.268.1. Migration to Version 8.0.12

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.268.2. Changes

• Remove overly-restrictive check for type length in constraints and functional indexes (Tom)

E.269. Release 8.0.11

Release date: 2007-02-05

This release contains a variety of fixes from 8.0.10, including a security fix. For information about new features in the 8.0 major release, see Section E.280.

E.269.1. Migration to Version 8.0.11

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.269.2. Changes

• Remove security vulnerabilities that allowed connected users to read backend memory (Tom)
The vulnerabilities involve suppressing the normal check that a SQL function returns the data type it's declared to, and changing the data type of a table column (CVE-2007-0555, CVE-2007-0556). These errors can easily be exploited to cause a backend crash, and in principle might be used to read database content that the user should not be able to access.

- Fix rare bug wherein btree index page splits could fail due to choosing an infeasible split point (Heikki Linnakangas)
- Fix for rare Assert() crash triggered by UNION (Tom)
- Tighten security of multi-byte character processing for UTF8 sequences over three bytes long (Tom)

### E.270. Release 8.0.10

**Release date:** 2007-01-08

This release contains a variety of fixes from 8.0.9. For information about new features in the 8.0 major release, see Section E.280.

#### E.270.1. Migration to Version 8.0.10

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

#### E.270.2. Changes

- Improve handling of `getaddinfo()` on AIX (Tom)
  
  This fixes a problem with starting the statistics collector, among other things.
- Fix “failed to re-find parent key” errors in VACUUM (Tom)
- Fix race condition for truncation of a large relation across a gigabyte boundary by VACUUM (Tom)
- Fix bugs affecting multi-gigabyte hash indexes (Tom)
- Fix possible deadlock in Windows signal handling (Teodor)
- Fix error when constructing an `ARRAY[]` made up of multiple empty elements (Tom)
- Fix ecpg memory leak during connection (Michael)
- `to_number()` and `to_char(numeric)` are now STABLE, not IMMUTABLE, for new initdb installs (Tom)

  This is because `lc_numeric` can potentially change the output of these functions.
- Improve index usage of regular expressions that use parentheses (Tom)

  This improves `psql \d` performance also.
- Update timezone database

This affects Australian and Canadian daylight-savings rules in particular.

### E.271. Release 8.0.9

**Release date:** 2006-10-16
This release contains a variety of fixes from 8.0.8. For information about new features in the 8.0 major release, see Section E.280.

E.271.1. Migration to Version 8.0.9

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

E.271.2. Changes

• Fix crash when referencing NEW row values in rule WHERE expressions (Tom)
• Fix core dump when an untyped literal is taken as ANYARRAY
• Fix mishandling of AFTER triggers when query contains a SQL function returning multiple rows (Tom)
• Fix ALTER TABLE ... TYPE to recheck NOT NULL for USING clause (Tom)
• Fix string_to_array() to handle overlapping matches for the separator string
  For example, string_to_array('123xx456xxx789', 'xx').
• Fix corner cases in pattern matching for psql's \d commands
• Fix index-corrupting bugs in /contrib/ltree (Teodor)
• Numerous robustness fixes in ecpg (Joachim Wieland)
• Fix backslash escaping in /contrib/dbmirror
• Fix instability of statistics collection on Win32 (Tom, Andrew)
• Fixes for AIX and Intel compilers (Tom)

E.272. Release 8.0.8

Release date: 2006-05-23

This release contains a variety of fixes from 8.0.7, including patches for extremely serious security issues. For information about new features in the 8.0 major release, see Section E.280.

E.272.1. Migration to Version 8.0.8

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.

Full security against the SQL-injection attacks described in CVE-2006-2313 and CVE-2006-2314 might require changes in application code. If you have applications that embed untrustworthy strings into SQL commands, you should examine them as soon as possible to ensure that they are using recommended escaping techniques. In most cases, applications should be using subroutines provided by libraries or drivers (such as libpq's PQescapeStringConn()) to perform string escaping, rather than relying on ad hoc code to do it.

E.272.2. Changes

• Change the server to reject invalidly-encoded multibyte characters in all cases (Tatsuo, Tom)
While PostgreSQL has been moving in this direction for some time, the checks are now applied uniformly to all encodings and all textual input, and are now always errors not merely warnings. This change defends against SQL-injection attacks of the type described in CVE-2006-2313.

• Reject unsafe uses of \ in string literals

As a server-side defense against SQL-injection attacks of the type described in CVE-2006-2314, the server now only accepts ' and not \ as a representation of ASCII single quote in SQL string literals. By default, \ is rejected only when client_encoding is set to a client-only encoding (SJIS, BIG5, GBK, GB18030, or UHC), which is the scenario in which SQL injection is possible. A new configuration parameter backslash_quote is available to adjust this behavior when needed. Note that full security against CVE-2006-2314 might require client-side changes; the purpose of backslash_quote is in part to make it obvious that insecure clients are insecure.

• Modify libpq's string-escaping routines to be aware of encoding considerations and standard_conforming_strings

This fixes libpq-using applications for the security issues described in CVE-2006-2313 and CVE-2006-2314, and also future-proofs them against the planned changeover to SQL-standard string literal syntax. Applications that use multiple PostgreSQL connections concurrently should migrate to PQescapeStringConn() and PQescapeByteaConn() to ensure that escaping is done correctly for the settings in use in each database connection. Applications that do string escaping “by hand” should be modified to rely on library routines instead.

• Fix some incorrect encoding conversion functions

  win1251_to_iso, alt_to_iso, euc_tw_to_big5, euc_tw_to_mic, mic_to_euc_tw were all broken to varying extents.

• Clean up stray remaining uses of \ in strings (Bruce, Jan)

• Fix bug that sometimes caused OR'd index scans to miss rows they should have returned

• Fix WAL replay for case where a btree index has been truncated

• Fix SIMILAR TO for patterns involving | (Tom)

• Fix SELECT INTO and CREATE TABLE AS to create tables in the default tablespace, not the base directory (Kris Jurka)

• Fix server to use custom DH SSL parameters correctly (Michael Fuhr)

• Fix for Bonjour on Intel Macs (Ashley Clark)

• Fix various minor memory leaks

• Fix problem with password prompting on some Win32 systems (Robert Kinberg)

**E.273. Release 8.0.7**

**Release date:** 2006-02-14

This release contains a variety of fixes from 8.0.6. For information about new features in the 8.0 major release, see Section E.280.

**E.273.1. Migration to Version 8.0.7**

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.6, see Section E.274.
E.273.2. Changes

- Fix potential crash in `SET SESSION AUTHORIZATION` (CVE-2006-0553)
  
  An unprivileged user could crash the server process, resulting in momentary denial of service to other users, if the server has been compiled with Asserts enabled (which is not the default). Thanks to Akio Ishida for reporting this problem.

- Fix bug with row visibility logic in self-inserted rows (Tom)
  
  Under rare circumstances a row inserted by the current command could be seen as already valid, when it should not be. Repairs bug created in 8.0.4, 7.4.9, and 7.3.11 releases.

- Fix race condition that could lead to “file already exists” errors during `pg_clog` and `pg_subtrans` file creation (Tom)

- Fix cases that could lead to crashes if a cache-invalidation message arrives at just the wrong time (Tom)

- Properly check `DOMAIN` constraints for `UNKNOWN` parameters in prepared statements (Neil)

- Ensure `ALTER COLUMN TYPE` will process `FOREIGN KEY`, `UNIQUE`, and `PRIMARY KEY` constraints in the proper order (Nakano Yoshihisa)

- Fixes to allow restoring dumps that have cross-schema references to custom operators or operator classes (Tom)

- Allow `pg_restore` to continue properly after a `COPY` failure; formerly it tried to treat the remaining `COPY` data as SQL commands (Stephen Frost)

- Fix `pg_ctl unregister` crash when the data directory is not specified (Magnus)

- Fix `ecpg` crash on AMD64 and PPC (Neil)

- Recover properly if error occurs during argument passing in PL/Python (Neil)

- Fix PL/Perl's handling of locales on Win32 to match the backend (Andrew)

- Fix crash when `log_min_messages` is set to `DEBUG3` or above in `postgresql.conf` on Win32 (Bruce)

- Fix `pgxs -L` library path specification for Win32, Cygwin, macOS, AIX (Bruce)

- Check that SID is enabled while checking for Win32 admin privileges (Magnus)

- Properly reject out-of-range date inputs (Kris Jurka)

- Portability fix for testing presence of `finite` and `isinf` during configure (Tom)

E.274. Release 8.0.6

Release date:  2006-01-09

This release contains a variety of fixes from 8.0.5. For information about new features in the 8.0 major release, see Section E.280.

E.274.1. Migration to Version 8.0.6

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.3, see Section E.277. Also, you might need to `REINDEX` indexes on textual columns after updating, if you are affected by the locale or plperl issues described below.
E.274.2. Changes

- Fix Windows code so that postmaster will continue rather than exit if there is no more room in ShmemBackendArray (Magnus)

  The previous behavior could lead to a denial-of-service situation if too many connection requests arrive close together. This applies only to the Windows port.

- Fix bug introduced in 8.0 that could allow ReadBuffer to return an already-used page as new, potentially causing loss of recently-committed data (Tom)

- Fix for protocol-level Describe messages issued outside a transaction or in a failed transaction (Tom)

- Fix character string comparison for locales that consider different character combinations as equal, such as Hungarian (Tom)

  This might require REINDEX to fix existing indexes on textual columns.

- Set locale environment variables during postmaster startup to ensure that plperl won't change the locale later

  This fixes a problem that occurred if the postmaster was started with environment variables specifying a different locale than what initdb had been told. Under these conditions, any use of plperl was likely to lead to corrupt indexes. You might need REINDEX to fix existing indexes on textual columns if this has happened to you.

- Allow more flexible relocation of installation directories (Tom)

  Previous releases supported relocation only if all installation directory paths were the same except for the last component.

- Fix longstanding bug in strpos() and regular expression handling in certain rarely used Asian multi-byte character sets (Tatsuo)

- Various fixes for functions returning RECORDs (Tom)

- Fix bug in /contrib/pgcrypto gen_salt, which caused it not to use all available salt space for MD5 and XDES algorithms (Marko Kreen, Solar Designer)

  Salts for Blowfish and standard DES are unaffected.

- Fix /contrib/dblink to throw an error, rather than crashing, when the number of columns specified is different from what's actually returned by the query (Joe)

E.275. Release 8.0.5

Release date: 2005-12-12

This release contains a variety of fixes from 8.0.4. For information about new features in the 8.0 major release, see Section E.280.

E.275.1. Migration to Version 8.0.5

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.3, see Section E.277.

E.275.2. Changes

- Fix race condition in transaction log management
There was a narrow window in which an I/O operation could be initiated for the wrong page, leading to an Assert failure or data corruption.

• Fix bgwriter problems after recovering from errors (Tom)

The background writer was found to leak buffer pins after write errors. While not fatal in itself, this might lead to mysterious blockages of later VACUUM commands.

• Prevent failure if client sends Bind protocol message when current transaction is already aborted

• /contrib/ltree fixes (Teodor)

• AIX and HPUX compile fixes (Tom)

• Retry file reads and writes after Windows NO_SYSTEM_RESOURCES error (Qingqing Zhou)

• Fix intermittent failure when log_line_prefix includes %i

• Fix psql performance issue with long scripts on Windows (Merlin Moncure)

• Fix missing updates of pg_group flat file

• Fix longstanding planning error for outer joins

  This bug sometimes caused a bogus error “RIGHT JOIN is only supported with merge-joinable join conditions”.

• Postpone timezone initialization until after postmaster.pid is created

  This avoids confusing startup scripts that expect the pid file to appear quickly.

• Prevent core dump in pg_autovacuum when a table has been dropped

• Fix problems with whole-row references (foo.*) to subquery results

E.276. Release 8.0.4

Release date: 2005-10-04

This release contains a variety of fixes from 8.0.3. For information about new features in the 8.0 major release, see Section E.280.

E.276.1. Migration to Version 8.0.4

A dump/restore is not required for those running 8.0.X. However, if you are upgrading from a version earlier than 8.0.3, see Section E.277.

E.276.2. Changes

• Fix error that allowed VACUUM to remove ctid chains too soon, and add more checking in code that follows ctid links

  This fixes a long-standing problem that could cause crashes in very rare circumstances.

• Fix CHAR() to properly pad spaces to the specified length when using a multiple-byte character set (Yoshiyuki Asaba)

  In prior releases, the padding of CHAR() was incorrect because it only padded to the specified number of bytes without considering how many characters were stored.
• Force a checkpoint before committing `CREATE DATABASE`.
  This should fix recent reports of “index is not a btree” failures when a crash occurs shortly after `CREATE DATABASE`.

• Fix the sense of the test for read-only transaction in `COPY`.
  The code formerly prohibited `COPY TO`, where it should prohibit `COPY FROM`.

• Handle consecutive embedded newlines in `COPY` CSV-mode input.

• Fix `date_trunc(week)` for dates near year end.

• Fix planning problem with outer-join ON clauses that reference only the inner-side relation.

• Further fixes for `x FULL JOIN y ON true` corner cases.

• Fix overenthusiastic optimization of `x IN (SELECT DISTINCT ...) and related cases`.

• Fix mis-planning of queries with small `LIMIT` values due to poorly thought out “fuzzy” cost comparison.

• Make `array_in` and `array_recv` more paranoid about validating their OID parameter.

• Fix missing rows in queries like `UPDATE a=... WHERE a... with GiST index on column a`.

• Improve robustness of datetime parsing.

• Improve checking for partially-written WAL pages.

• Improve robustness of signal handling when SSL is enabled.

• Improve MIPS and M68K spinlock code.

• Don’t try to open more than `max_files_per_process` files during postmaster startup.

• Various memory leakage fixes.

• Various portability improvements.

• Update timezone data files.

• Improve handling of DLL load failures on Windows.

• Improve random-number generation on Windows.

• Make `psql -f filename return a nonzero exit code when opening the file fails`.

• Change `pg_dump` to handle inherited check constraints more reliably.

• Fix password prompting in `pg_restore` on Windows.

• Fix PL/pgSQL to handle `var := var` correctly when the variable is of pass-by-reference type.

• Fix PL/Perl `%_SHARED` so it’s actually shared.

• Fix `contrib/pg_autovacuum` to allow sleep intervals over 2000 sec.

• Update `contrib/tsearch2` to use current Snowball code.

### E.277. Release 8.0.3

**Release date:** 2005-05-09
This release contains a variety of fixes from 8.0.2, including several security-related issues. For information about new features in the 8.0 major release, see Section E.280.

E.277.1. Migration to Version 8.0.3

A dump/restore is not required for those running 8.0.X. However, it is one possible way of handling two significant security problems that have been found in the initial contents of 8.0.X system catalogs. A dump/initdb/reload sequence using 8.0.3's initdb will automatically correct these problems.

The larger security problem is that the built-in character set encoding conversion functions can be invoked from SQL commands by unprivileged users, but the functions were not designed for such use and are not secure against malicious choices of arguments. The fix involves changing the declared parameter list of these functions so that they can no longer be invoked from SQL commands. (This does not affect their normal use by the encoding conversion machinery.)

The lesser problem is that the contrib/tsearch2 module creates several functions that are improperly declared to return internal when they do not accept internal arguments. This breaks type safety for all functions using internal arguments.

It is strongly recommended that all installations repair these errors, either by initdb or by following the manual repair procedure given below. The errors at least allow unprivileged database users to crash their server process, and might allow unprivileged users to gain the privileges of a database superuser.

If you wish not to do an initdb, perform the same manual repair procedures shown in the 7.4.8 release notes.

E.277.2. Changes

• Change encoding function signature to prevent misuse
• Change contrib/tsearch2 to avoid unsafe use of INTERNAL function results
• Guard against incorrect second parameter to record_out
• Repair ancient race condition that allowed a transaction to be seen as committed for some purposes (eg SELECT FOR UPDATE) slightly sooner than for other purposes

This is an extremely serious bug since it could lead to apparent data inconsistencies being briefly visible to applications.

• Repair race condition between relation extension and VACUUM

This could theoretically have caused loss of a page's worth of freshly-inserted data, although the scenario seems of very low probability. There are no known cases of it having caused more than an Assert failure.

• Fix comparisons of TIME WITH TIME ZONE values

The comparison code was wrong in the case where the --enable-integer-datetimes configuration switch had been used. NOTE: if you have an index on a TIME WITH TIME ZONE column, it will need to be REINDEXed after installing this update, because the fix corrects the sort order of column values.

• Fix EXTRACT(EPOCH) for TIME WITH TIME ZONE values

• Fix mis-display of negative fractional seconds in INTERVAL values

This error only occurred when the --enable-integer-datetimes configuration switch had been used.

• Fix pg_dump to dump trigger names containing % correctly (Neil)
Release Notes

• Still more 64-bit fixes for contrib/intagg
• Prevent incorrect optimization of functions returning RECORD
• Prevent crash on COALESCE (NULL, NULL)
• Fix Borland makefile for libpq
• Fix contrib/btree_gist for timetz type (Teodor)
• Make pg_ctl check the PID found in postmaster.pid to see if it is still a live process
• Fix pg_dump/pg_restore problems caused by addition of dump timestamps
• Fix interaction between materializing holdable cursors and firing deferred triggers during transaction commit
• Fix memory leak in SQL functions returning pass-by-reference data types

E.278. Release 8.0.2

Release date: 2005-04-07

This release contains a variety of fixes from 8.0.1. For information about new features in the 8.0 major release, see Section E.280.

E.278.1. Migration to Version 8.0.2

A dump/restore is not required for those running 8.0.*. This release updates the major version number of the PostgreSQL libraries, so it might be necessary to re-link some user applications if they cannot find the properly-numbered shared library.

E.278.2. Changes

• Increment the major version number of all interface libraries (Bruce)

This should have been done in 8.0.0. It is required so 7.4.X versions of PostgreSQL client applications, like psql, can be used on the same machine as 8.0.X applications. This might require re-linking user applications that use these libraries.

• Add Windows-only wal_sync_method setting of fsync_writethrough (Magnus, Bruce)

This setting causes PostgreSQL to write through any disk-drive write cache when writing to WAL. This behavior was formerly called fsync, but was renamed because it acts quite differently from fsync on other platforms.

• Enable the wal_sync_method setting of open_datasync on Windows, and make it the default for that platform (Magnus, Bruce)

Because the default is no longer fsync_writethrough, data loss is possible during a power failure if the disk drive has write caching enabled. To turn off the write cache on Windows, from the Device Manager, choose the drive properties, then Policies.

• New cache management algorithm 2Q replaces ARC (Tom)

This was done to avoid a pending US patent on ARC. The 2Q code might be a few percentage points slower than ARC for some work loads. A better cache management algorithm will appear in 8.1.

• Planner adjustments to improve behavior on freshly-created tables (Tom)
• Allow plpgsql to assign to an element of an array that is initially NULL (Tom)

Formerly the array would remain NULL, but now it becomes a single-element array. The main SQL engine was changed to handle UPDATE of a null array value this way in 8.0, but the similar case in plpgsql was overlooked.

• Convert \r\n and \r to \n in plpython function bodies (Michael Fuhr)

This prevents syntax errors when plpython code is written on a Windows or Mac client.

• Allow SPI cursors to handle utility commands that return rows, such as EXPLAIN (Tom)

• Fix CLUSTER failure after ALTER TABLE SET WITHOUT OIDS (Tom)

• Reduce memory usage of ALTER TABLE ADD COLUMN (Neil)

• Fix ALTER LANGUAGE RENAME (Tom)

• Document the Windows-only register and unregister options of pg_ctl (Magnus)

• Ensure operations done during backend shutdown are counted by statistics collector

This is expected to resolve reports of pg_autovacuum not vacuuming the system catalogs often enough — it was not being told about catalog deletions caused by temporary table removal during backend exit.

• Change the Windows default for configuration parameter log_destination to eventlog (Magnus)

By default, a server running on Windows will now send log output to the Windows event logger rather than standard error.

• Make Kerberos authentication work on Windows (Magnus)

• Allow ALTER DATABASE RENAME by superusers who aren’t flagged as having CREATEDB privilege (Tom)

• Modify WAL log entries for CREATE and DROP DATABASE to not specify absolute paths (Tom)

This allows point-in-time recovery on a different machine with possibly different database location. Note that CREATE TABLESPACE still poses a hazard in such situations.

• Fix crash from a backend exiting with an open transaction that created a table and opened a cursor on it (Tom)

• Fix array_map() so it can call PL functions (Tom)

• Several contrib/tsearch2 and contrib/btree_gist fixes (Teodor)

• Fix crash of some contrib/pgcrypto functions on some platforms (Marko Kreen)

• Fix contrib/intagg for 64-bit platforms (Tom)

• Fix ecpg bugs in parsing of CREATE statement (Michael)

• Work around gcc bug on powerpc and amd64 causing problems in ecpg (Christof Petig)

• Do not use locale-aware versions of upper(), lower(), and initcap() when the locale is C (Bruce)

This allows these functions to work on platforms that generate errors for non-7-bit data when the locale is C.
• Fix \texttt{quote_ident()} to quote names that match keywords (Tom)

• Fix \texttt{to_date()} to behave reasonably when \texttt{CC} and \texttt{YY} fields are both used (Karel)

• Prevent \texttt{to_char(interval)} from failing when given a zero-month interval (Tom)

• Fix wrong week returned by \texttt{date_trunc('week')} (Bruce)

\texttt{date_trunc('week')} returned the wrong year for the first few days of January in some years.

• Use the correct default mask length for class \texttt{D} addresses in \texttt{INET} data types (Tom)

\section*{E.279. Release 8.0.1}

Release date: 2005-01-31

This release contains a variety of fixes from 8.0.0, including several security-related issues. For information about new features in the 8.0 major release, see Section E.280.

\subsection*{E.279.1. Migration to Version 8.0.1}

A dump/restore is not required for those running 8.0.0.

\subsection*{E.279.2. Changes}

• Disallow \texttt{LOAD} to non-superusers

On platforms that will automatically execute initialization functions of a shared library (this includes at least Windows and ELF-based Unixen), \texttt{LOAD} can be used to make the server execute arbitrary code. Thanks to NGS Software for reporting this.

• Check that creator of an aggregate function has the right to execute the specified transition functions

This oversight made it possible to bypass denial of \texttt{EXECUTE} permission on a function.

• Fix security and 64-bit issues in contrib/intagg

• Add needed \texttt{STRICT} marking to some contrib functions (Kris Jurka)

• Avoid buffer overrun when \texttt{plpgsql} cursor declaration has too many parameters (Neil)

• Make \texttt{ALTER TABLE ADD COLUMN} enforce domain constraints in all cases

• Fix planning error for \texttt{FULL} and \texttt{RIGHT} outer joins

The result of the join was mistakenly supposed to be sorted the same as the left input. This could not only deliver mis-sorted output to the user, but in case of nested merge joins could give outright wrong answers.

• Improve planning of grouped aggregate queries

• \texttt{ROLLBACK TO savepoint} closes cursors created since the savepoint

• Fix inadequate backend stack size on Windows

• Avoid \texttt{SHGetSpecialFolderPath()} on Windows (Magnus)

• Fix some problems in running \texttt{pg_autovacuum} as a Windows service (Dave Page)

• Multiple minor bug fixes in \texttt{pg_dump/pg_restore}
E.280. Release 8.0

Release date: 2005-01-19

E.280.1. Overview

Major changes in this release:

Microsoft Windows Native Server

This is the first PostgreSQL release to run natively on Microsoft Windows® as a server. It can run as a Windows service. This release supports NT-based Windows releases like Windows 2000 SP4, Windows XP, and Windows 2003. Older releases like Windows 95, Windows 98, and Windows ME are not supported because these operating systems do not have the infrastructure to support PostgreSQL. A separate installer project has been created to ease installation on Windows — see http://www.postgresql.org/ftp/win32/.

Although tested throughout our release cycle, the Windows port does not have the benefit of years of use in production environments that PostgreSQL has on Unix platforms. Therefore it should be treated with the same level of caution as you would a new product.

Previous releases required the Unix emulation toolkit Cygwin in order to run the server on Windows operating systems. PostgreSQL has supported native clients on Windows for many years.

Savepoints

Savepoints allow specific parts of a transaction to be aborted without affecting the remainder of the transaction. Prior releases had no such capability; there was no way to recover from a statement failure within a transaction except by aborting the whole transaction. This feature is valuable for application writers who require error recovery within a complex transaction.

Point-In-Time Recovery

In previous releases there was no way to recover from disk drive failure except to restore from a previous backup or use a standby replication server. Point-in-time recovery allows continuous backup of the server. You can recover either to the point of failure or to some transaction in the past.

Tablespaces

Tablespaces allow administrators to select different file systems for storage of individual tables, indexes, and databases. This improves performance and control over disk space usage. Prior releases used initlocation and manual symlink management for such tasks.

Improved Buffer Management, CHECKPOINT, VACUUM

This release has a more intelligent buffer replacement strategy, which will make better use of available shared buffers and improve performance. The performance impact of vacuum and checkpoints is also lessened.

Change Column Types

A column's data type can now be changed with ALTER TABLE.

New Perl Server-Side Language

A new version of the plperl server-side language now supports a persistent shared storage area, triggers, returning records and arrays of records, and SPI calls to access the database.
Comma-separated-value (CSV) support in **COPY**

**COPY** can now read and write comma-separated-value files. It has the flexibility to interpret non-standard quoting and separation characters too.

### E.280.2. Migration to Version 8.0

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

- In **READ COMMITTED** serialization mode, volatile functions now see the results of concurrent transactions committed up to the beginning of each statement within the function, rather than up to the beginning of the interactive command that called the function.

- Functions declared **STABLE** or **IMMUTABLE** always use the snapshot of the calling query, and therefore do not see the effects of actions taken after the calling query starts, whether in their own transaction or other transactions. Such a function must be read-only, too, meaning that it cannot use any SQL commands other than **SELECT**.

- Nondeferred **AFTER** triggers are now fired immediately after completion of the triggering query, rather than upon finishing the current interactive command. This makes a difference when the triggering query occurred within a function: the trigger is invoked before the function proceeds to its next operation.

- Server configuration parameters **virtual_host** and **tcip_socket** have been replaced with a more general parameter **listen_addresses**. Also, the server now listens on **localhost** by default, which eliminates the need for the `-i` postmaster switch in many scenarios.

- Server configuration parameters **SortMem** and **VacuumMem** have been renamed to **work_mem** and **maintenance_work_mem** to better reflect their use. The original names are still supported in **SET** and **SHOW**.

- Server configuration parameters **log_pid**, **log_timestamp**, and **log_source_port** have been replaced with a more general parameter **log_line_prefix**.

- Server configuration parameter **syslog** has been replaced with a more logical **log_destination** variable to control the log output destination.

- Server configuration parameter **log_statement** has been changed so it can selectively log just database modification or data definition statements. Server configuration parameter **log_duration** now prints only when **log_statement** prints the query.

- Server configuration parameter **max_expr_depth** parameter has been replaced with **max_stack_depth** which measures the physical stack size rather than the expression nesting depth. This helps prevent session termination due to stack overflow caused by recursive functions.

- The **length()** function no longer counts trailing spaces in **CHAR(n)** values.

- Casting an integer to **BIT(N)** selects the rightmost **N** bits of the integer, not the leftmost **N** bits as before.

- Updating an element or slice of a NULL array value now produces a nonnull array result, namely an array containing just the assigned-to positions.

- Syntax checking of array input values has been tightened up considerably. Junk that was previously allowed in odd places with odd results now causes an error. Empty-string element values must now be written as " ", rather than writing nothing. Also changed behavior with respect to whitespace surrounding array elements: trailing whitespace is now ignored, for symmetry with leading whitespace (which has always been ignored).

- Overflow in integer arithmetic operations is now detected and reported as an error.
• The arithmetic operators associated with the single-byte "char" data type have been removed.

• The extract() function (also called date_part) now returns the proper year for BC dates. It previously returned one less than the correct year. The function now also returns the proper values for millennium and century.

• CIDR values now must have their nonmasked bits be zero. For example, we no longer allow 204.248.199.1/31 as a CIDR value. Such values should never have been accepted by PostgreSQL and will now be rejected.

• EXECUTE now returns a completion tag that matches the executed statement.

• psql's \copy command now reads or writes to the query's stdin/stdout, rather than psql's stdin/stdout. The previous behavior can be accessed via new pstdin/pstdout parameters.

• The JDBC client interface has been removed from the core distribution, and is now hosted at http://jdbc.postgresql.org.

• The Tcl client interface has also been removed. There are several Tcl interfaces now hosted at http://gborg.postgresql.org.

• The server now uses its own time zone database, rather than the one supplied by the operating system. This will provide consistent behavior across all platforms. In most cases, there should be little noticeable difference in time zone behavior, except that the time zone names used by SET/SHOW TimeZone might be different from what your platform provides.

• Configure's threading option no longer requires users to run tests or edit configuration files; threading options are now detected automatically.

• Now that tablespaces have been implemented, initlocation has been removed.

• The API for user-defined GiST indexes has been changed. The Union and PickSplit methods are now passed a pointer to a special GistEntryVector structure, rather than a bytea.

E.280.3. Deprecated Features

Some aspects of PostgreSQL's behavior have been determined to be suboptimal. For the sake of backward compatibility these have not been removed in 8.0, but they are considered deprecated and will be removed in the next major release.

• The 8.1 release will remove the to_char() function for intervals.

• The server now warns of empty strings passed to oid/float4/float8 data types, but continues to interpret them as zeroes as before. In the next major release, empty strings will be considered invalid input for these data types.

• By default, tables in PostgreSQL 8.0 and earlier are created with OIDs. In the next release, this will not be the case: to create a table that contains OIDs, the WITH OIDS clause must be specified or the default_with_oids configuration parameter must be set. Users are encouraged to explicitly specify WITH OIDS if their tables require OIDs for compatibility with future releases of PostgreSQL.

E.280.4. Changes

Below you will find a detailed account of the changes between release 8.0 and the previous major release.

E.280.4.1. Performance Improvements

• Support cross-data-type index usage (Tom)
Before this change, many queries would not use an index if the data types did not match exactly. This improvement makes index usage more intuitive and consistent.

- New buffer replacement strategy that improves caching (Jan)

Prior releases used a least-recently-used (LRU) cache to keep recently referenced pages in memory. The LRU algorithm did not consider the number of times a specific cache entry was accessed, so large table scans could force out useful cache pages. The new cache algorithm uses four separate lists to track most recently used and most frequently used cache pages and dynamically optimize their replacement based on the work load. This should lead to much more efficient use of the shared buffer cache. Administrators who have tested shared buffer sizes in the past should retest with this new cache replacement policy.

- Add subprocess to write dirty buffers periodically to reduce checkpoint writes (Jan)

In previous releases, the checkpoint process, which runs every few minutes, would write all dirty buffers to the operating system's buffer cache then flush all dirty operating system buffers to disk. This resulted in a periodic spike in disk usage that often hurt performance. The new code uses a background writer to trickle disk writes at a steady pace so checkpoints have far fewer dirty pages to write to disk. Also, the new code does not issue a global sync() call, but instead fsync()s just the files written since the last checkpoint. This should improve performance and minimize degradation during checkpoints.

- Add ability to prolong vacuum to reduce performance impact (Jan)

On busy systems, VACUUM performs many I/O requests which can hurt performance for other users. This release allows you to slow down VACUUM to reduce its impact on other users, though this increases the total duration of VACUUM.

- Improve B-tree index performance for duplicate keys (Dmitry Tkach, Tom)

This improves the way indexes are scanned when many duplicate values exist in the index.

- Use dynamically-generated table size estimates while planning (Tom)

Formerly the planner estimated table sizes using the values seen by the last VACUUM or ANALYZE, both as to physical table size (number of pages) and number of rows. Now, the current physical table size is obtained from the kernel, and the number of rows is estimated by multiplying the table size by the row density (rows per page) seen by the last VACUUM or ANALYZE. This should produce more reliable estimates in cases where the table size has changed significantly since the last housekeeping command.

- Improved index usage with OR clauses (Tom)

This allows the optimizer to use indexes in statements with many OR clauses that would not have been indexed in the past. It can also use multi-column indexes where the first column is specified and the second column is part of an OR clause.

- Improve matching of partial index clauses (Tom)

The server is now smarter about using partial indexes in queries involving complex WHERE clauses.

- Improve performance of the GEQO optimizer (Tom)

The GEQO optimizer is used to plan queries involving many tables (by default, twelve or more). This release speeds up the way queries are analyzed to decrease time spent in optimization.

- Miscellaneous optimizer improvements

There is not room here to list all the minor improvements made, but numerous special cases work better than in prior releases.
• Improve lookup speed for C functions (Tom)

This release uses a hash table to lookup information for dynamically loaded C functions. This improves their speed so they perform nearly as quickly as functions that are built into the server executable.

• Add type-specific ANALYZE statistics capability (Mark Cave-Ayland)

This feature allows more flexibility in generating statistics for nonstandard data types.

• ANALYZE now collects statistics for expression indexes (Tom)

Expression indexes (also called functional indexes) allow users to index not just columns but the results of expressions and function calls. With this release, the optimizer can gather and use statistics about the contents of expression indexes. This will greatly improve the quality of planning for queries in which an expression index is relevant.

• New two-stage sampling method for ANALYZE (Manfred Koizar)

This gives better statistics when the density of valid rows is very different in different regions of a table.

• Speed up TRUNCATE (Tom)

This buys back some of the performance loss observed in 7.4, while still keeping TRUNCATE transaction-safe.

E.280.4.2. Server Changes

• Add WAL file archiving and point-in-time recovery (Simon Riggs)

• Add tablespaces so admins can control disk layout (Gavin)

• Add a built-in log rotation program (Andreas Pflug)

It is now possible to log server messages conveniently without relying on either syslog or an external log rotation program.

• Add new read-only server configuration parameters to show server compile-time settings: block_size, integer_datetimes, max_function_args, max_identifier_length, max_index_keys (Joe)

• Make quoting of sameuser, samegroup, and all remove special meaning of these terms in pg_hba.conf (Andrew)

• Use clearer IPv6 name ::1/128 for localhost in default pg_hba.conf (Andrew)

• Use CIDR format in pg_hba.conf examples (Andrew)

• Rename server configuration parameters SortMem and VacuumMem to work_mem and maintenance_work_mem (Old names still supported) (Tom)

This change was made to clarify that bulk operations such as index and foreign key creation use maintenance_work_mem, while work_mem is for workspaces used during query execution.

• Allow logging of session disconnections using server configuration log_disconnections (Andrew)

• Add new server configuration parameter log_line_prefix to allow control of information emitted in each log line (Andrew)

Available information includes user name, database name, remote IP address, and session start time.
• Remove server configuration parameters log_pid, log_timestamp, log_source_port; functionality superseded by log_line_prefix (Andrew)

• Replace the virtual_host and tcpip_socket parameters with a unified listen_addresses parameter (Andrew, Tom)

  virtual_host could only specify a single IP address to listen on. listen_addresses allows multiple addresses to be specified.

• Listen on localhost by default, which eliminates the need for the -i postmaster switch in many scenarios (Andrew)

  Listening on localhost (127.0.0.1) opens no new security holes but allows configurations like Windows and JDBC, which do not support local sockets, to work without special adjustments.

• Remove syslog server configuration parameter, and add more logical log_destination variable to control log output location (Magnus)

• Change server configuration parameter log_statement to take values all, mod, ddl, or none to select which queries are logged (Bruce)

  This allows administrators to log only data definition changes or only data modification statements.

• Some logging-related configuration parameters could formerly be adjusted by ordinary users, but only in the “more verbose” direction. They are now treated more strictly: only superusers can set them. However, a superuser can use ALTER USER to provide per-user settings of these values for non-superusers. Also, it is now possible for superusers to set values of superuser-only configuration parameters via PGOPTIONS.

• Allow configuration files to be placed outside the data directory (mlw)

  By default, configuration files are kept in the cluster's top directory. With this addition, configuration files can be placed outside the data directory, easing administration.

• Plan prepared queries only when first executed so constants can be used for statistics (Oliver Jowett)

  Prepared statements plan queries once and execute them many times. While prepared queries avoid the overhead of re-planning on each use, the quality of the plan suffers from not knowing the exact parameters to be used in the query. In this release, planning of unnamed prepared statements is delayed until the first execution, and the actual parameter values of that execution are used as optimization hints. This allows use of out-of-line parameter passing without incurring a performance penalty.

• Allow DECLARE CURSOR to take parameters (Oliver Jowett)

  It is now useful to issue DECLARE CURSOR in a Parse message with parameters. The parameter values sent at Bind time will be substituted into the execution of the cursor's query.

• Fix hash joins and aggregates of inet and cidr data types (Tom)

  Release 7.4 handled hashing of mixed inet and cidr values incorrectly. (This bug did not exist in prior releases because they wouldn't try to hash either data type.)

• Make log_duration print only when log_statement prints the query (Ed L.)

E.280.4.3. Query Changes

• Add savepoints (nested transactions) (Alvaro)

• Unsupported isolation levels are now accepted and promoted to the nearest supported level (Peter)
The SQL specification states that if a database doesn't support a specific isolation level, it should use the next more restrictive level. This change complies with that recommendation.

- Allow `BEGIN WORK` to specify transaction isolation levels like `START TRANSACTION` does (Bruce)

- Fix table permission checking for cases in which rules generate a query type different from the originally submitted query (Tom)

- Implement dollar quoting to simplify single-quote usage (Andrew, Tom, David Fetter)

  In previous releases, because single quotes had to be used to quote a function's body, the use of single quotes inside the function text required use of two single quotes or other error-prone notations. With this release we add the ability to use "dollar quoting" to quote a block of text. The ability to use different quoting delimiters at different nesting levels greatly simplifies the task of quoting correctly, especially in complex functions. Dollar quoting can be used anywhere quoted text is needed.

- Make `CASE val WHEN compval1 THEN ... evaluate val only once` (Tom)

  `CASE` no longer evaluates the tested expression multiple times. This has benefits when the expression is complex or is volatile.

- Test `HAVING` before computing target list of an aggregate query (Tom)

  Fixes improper failure of cases such as `SELECT SUM(win)/SUM(lose) ... GROUP BY ... HAVING SUM(lose) > 0`. This should work but formerly could fail with divide-by-zero.

- Replace `max_expr_depth` parameter with `max_stack_depth` parameter, measured in kilobytes of stack size (Tom)

  This gives us a fairly bulletproof defense against crashing due to runaway recursive functions. Instead of measuring the depth of expression nesting, we now directly measure the size of the execution stack.

- Allow arbitrary row expressions (Tom)

  This release allows SQL expressions to contain arbitrary composite types, that is, row values. It also allows functions to more easily take rows as arguments and return row values.

- Allow `LIKE/ILIKE` to be used as the operator in row and subselect comparisons (Fabien Coelho)

- Avoid locale-specific case conversion of basic ASCII letters in identifiers and keywords (Tom)

  This solves the “Turkish problem” with mangling of words containing İ and ı. Folding of characters outside the 7-bit-ASCII set is still locale-aware.

- Improve syntax error reporting (Fabien, Tom)

  Syntax error reports are more useful than before.

- Change `EXECUTE` to return a completion tag matching the executed statement (Kris Jurka)

  Previous releases return an `EXECUTE` tag for any `EXECUTE` call. In this release, the tag returned will reflect the command executed.

- Avoid emitting `NATURAL CROSS JOIN` in rule listings (Tom)

  Such a clause makes no logical sense, but in some cases the rule decompiler formerly produced this syntax.
E.280.4.4. Object Manipulation Changes

- Add `COMMENT` ON for casts, conversions, languages, operator classes, and large objects (Christopher)

- Add new server configuration parameter `default_with_oids` to control whether tables are created with OIDs by default (Neil)

  This allows administrators to control whether CREATE TABLE commands create tables with or without OID columns by default. (Note: the current factory default setting for `default_with_oids` is TRUE, but the default will become FALSE in future releases.)

- Add `WITH / WITHOUT OIDS` clause to CREATE TABLE AS (Neil)

- Allow `ALTER TABLE DROP COLUMN` to drop an OID column (`ALTER TABLE SET WITHOUT OIDS` still works) (Tom)

- Allow composite types as table columns (Tom)

- Allow `ALTER ... ADD COLUMN` with defaults and NOT NULL constraints; works per SQL spec (Rod)

  It is now possible for `ADD COLUMN` to create a column that is not initially filled with NULLs, but with a specified default value.

- Add `ALTER COLUMN TYPE` to change column's type (Rod)

  It is now possible to alter a column's data type without dropping and re-adding the column.

- Allow multiple `ALTER` actions in a single `ALTER TABLE` command (Rod)

  This is particularly useful for `ALTER` commands that rewrite the table (which include `ALTER COLUMN TYPE` and `ADD COLUMN` with a default). By grouping `ALTER` commands together, the table need be rewritten only once.

- Allow `ALTER TABLE` to add SERIAL columns (Tom)

  This falls out from the new capability of specifying defaults for new columns.

- Allow changing the owners of aggregates, conversions, databases, functions, operators, operator classes, schemas, types, and tablespaces (Christopher, Euler Taveira de Oliveira)

  Previously this required modifying the system tables directly.

- Allow temporary object creation to be limited to `SECURITY DEFINER` functions (Sean Chitten-den)

- Add `ALTER TABLE ... SET WITHOUT CLUSTER` (Christopher)

  Prior to this release, there was no way to clear an auto-cluster specification except to modify the system tables.

- Constraint/Index/SERIAL names are now `table_column_type` with numbers appended to guarantee uniqueness within the schema (Tom)

  The SQL specification states that such names should be unique within a schema.

- Add `pg_get_serial_sequence()` to return a SERIAL column's sequence name (Christopher)

  This allows automated scripts to reliably find the SERIAL sequence name.
• Warn when primary/foreign key data type mismatch requires costly lookup
• New ALTER INDEX command to allow moving of indexes between tablespaces (Gavin)
• Make ALTER TABLE OWNER change dependent sequence ownership too (Alvaro)

E.280.4.5. Utility Command Changes

• Allow CREATE SCHEMA to create triggers, indexes, and sequences (Neil)
• Add ALSO keyword to CREATE RULE (Fabien Coelho)
  This allows ALSO to be added to rule creation to contrast it with INSTEAD rules.
• Add NOWAIT option to LOCK (Tatsuo)
  This allows the LOCK command to fail if it would have to wait for the requested lock.
• Allow COPY to read and write comma-separated-value (CSV) files (Andrew, Bruce)
• Generate error if the COPY delimiter and NULL string conflict (Bruce)
• GRANT/REVOKE behavior follows the SQL spec more closely
• Avoid locking conflict between CREATE INDEX and CHECKPOINT (Tom)
  In 7.3 and 7.4, a long-running B-tree index build could block concurrent CHECKPOINTs from completing, thereby causing WAL bloat because the WAL log could not be recycled.
• Database-wide ANALYZE does not hold locks across tables (Tom)
  This reduces the potential for deadlocks against other backends that want exclusive locks on tables. To get the benefit of this change, do not execute database-wide ANALYZE inside a transaction block (BEGIN block); it must be able to commit and start a new transaction for each table.
• REINDEX does not exclusively lock the index's parent table anymore
  The index itself is still exclusively locked, but readers of the table can continue if they are not using the particular index being rebuilt.
• Erase MD5 user passwords when a user is renamed (Bruce)
  PostgreSQL uses the user name as salt when encrypting passwords via MD5. When a user's name is changed, the salt will no longer match the stored MD5 password, so the stored password becomes useless. In this release a notice is generated and the password is cleared. A new password must then be assigned if the user is to be able to log in with a password.
• New pg_ctl kill option for Windows (Andrew)
  Windows does not have a kill command to send signals to backends so this capability was added to pg_ctl.
• Information schema improvements
• Add --pwfile option to initdb so the initial password can be set by GUI tools (Magnus)
• Detect locale/encoding mismatch in initdb (Peter)
• Add register command to pg_ctl to register Windows operating system service (Dave Page)

E.280.4.6. Data Type and Function Changes

• More complete support for composite types (row types) (Tom)
Composite values can be used in many places where only scalar values worked before.

- Reject nonrectangular array values as erroneous (Joe)
  Formerly, array_in would silently build a surprising result.

- Overflow in integer arithmetic operations is now detected (Tom)

- The arithmetic operators associated with the single-byte "char" data type have been removed.
  Formerly, the parser would select these operators in many situations where an "unable to select an
  operator" error would be more appropriate, such as null * null. If you actually want to do
  arithmetic on a "char" column, you can cast it to integer explicitly.

- Syntax checking of array input values considerably tightened up (Joe)
  Junk that was previously allowed in odd places with odd results now causes an ERROR, for example,
  non-whitespace after the closing right brace.

- Empty-string array element values must now be written as "", rather than writing nothing (Joe)
  Formerly, both ways of writing an empty-string element value were allowed, but now a quoted
  empty string is required. The case where nothing at all appears will probably be considered to be
  a NULL element value in some future release.

- Array element trailing whitespace is now ignored (Joe)
  Formerly leading whitespace was ignored, but trailing whitespace between an element value and
  the delimiter or right brace was significant. Now trailing whitespace is also ignored.

- Emit array values with explicit array bounds when lower bound is not one (Joe)

- Accept YYYY-monthname-DD as a date string (Tom)

- Make netmask and hostmask functions return maximum-length mask length (Tom)

- Change factorial function to return numeric (Gavin)
  Returning numeric allows the factorial function to work for a wider range of input values.

- to_char/to_date() date conversion improvements (Kurt Roeckx, Fabien Coelho)

- Make length() disregard trailing spaces in CHAR(n) (Gavin)
  This change was made to improve consistency: trailing spaces are semantically insignificant in
  CHAR(n) data, so they should not be counted by length().

- Warn about empty string being passed to OID/float4/float8 data types (Neil)
  8.1 will throw an error instead.

- Allow leading or trailing whitespace in int2/int4/int8/float4/float8 input routines (Neil)

- Better support for IEEE Infinity and NaN values in float4/float8 (Neil)
  These should now work on all platforms that support IEEE-compliant floating point arithmetic.

- Add week option to date_trunc() (Robert Creager)

- Fix to_char for 1 BC (previously it returned 1 AD) (Bruce)

- Fix date_part(year) for BC dates (previously it returned one less than the correct year) (Bruce)
• Fix \texttt{date\_part()} to return the proper millennium and century (Fabien Coelho)

In previous versions, the century and millennium results had a wrong number and started in the
wrong year, as compared to standard reckoning of such things.

• Add \texttt{ceiling()} as an alias for \texttt{ceil()}, and \texttt{power()} as an alias for \texttt{pow()} for standards
compliance (Neil)

• Change \texttt{ln()}, \texttt{log()}, \texttt{power()}, and \texttt{sqrt()} to emit the correct SQLSTATE error codes for
certain error conditions, as specified by SQL:2003 (Neil)

• Add \texttt{width\_bucket()} function as defined by SQL:2003 (Neil)

• Add \texttt{generate\_series()} functions to simplify working with numeric sets (Joe)

• Fix upper/lower/initcap() functions to work with multibyte encodings (Tom)

• Add boolean and bitwise integer AND/OR aggregates (Fabien Coelho)

• New session information functions to return network addresses for client and server (Sean Chitten-
den)

• Add function to determine the area of a closed path (Sean Chittenden)

• Add function to send cancel request to other backends (Magnus)

• Add \texttt{interval} plus \texttt{datetime} operators (Tom)

The reverse ordering, \texttt{datetime} plus \texttt{interval}, was already supported, but both are required
by the SQL standard.

• Casting an integer to \texttt{BIT(N)} selects the rightmost N bits of the integer (Tom)

In prior releases, the leftmost N bits were selected, but this was deemed unhelpful, not to mention
inconsistent with casting from bit to int.

• Require CIDR values to have all nonmasked bits be zero (Kevin Brintnall)

\section*{E.280.4.7. Server-Side Language Changes}

• In \texttt{READ COMMITTED} serialization mode, volatile functions now see the results of concurrent
transactions committed up to the beginning of each statement within the function, rather than up to
the beginning of the interactive command that called the function.

• Functions declared \texttt{STABLE} or \texttt{IMMUTABLE} always use the snapshot of the calling query, and
therefore do not see the effects of actions taken after the calling query starts, whether in their own
transaction or other transactions. Such a function must be read-only, too, meaning that it cannot use
any SQL commands other than \texttt{SELECT}. There is a considerable performance gain from declaring
a function \texttt{STABLE} or \texttt{IMMUTABLE} rather than \texttt{VOLATILE}.

• Nondeferred \texttt{AFTER} triggers are now fired immediately after completion of the triggering query,
rather than upon finishing the current interactive command. This makes a difference when the trig-
gering query occurred within a function: the trigger is invoked before the function proceeds to its
next operation. For example, if a function inserts a new row into a table, any nondeferred foreign
key checks occur before proceeding with the function.

• Allow function parameters to be declared with names (Dennis Björklund)

This allows better documentation of functions. Whether the names actually do anything depends
on the specific function language being used.

• Allow PL/pgSQL parameter names to be referenced in the function (Dennis Björklund)
This basically creates an automatic alias for each named parameter.

- Do minimal syntax checking of PL/pgSQL functions at creation time (Tom)

  This allows us to catch simple syntax errors sooner.

- More support for composite types (row and record variables) in PL/pgSQL

  For example, it now works to pass a rowtype variable to another function as a single variable.

- Default values for PL/pgSQL variables can now reference previously declared variables

- Improve parsing of PL/pgSQL FOR loops (Tom)

  Parsing is now driven by presence of ".." rather than data type of FOR variable. This makes no difference for correct functions, but should result in more understandable error messages when a mistake is made.

- Major overhaul of PL/Perl server-side language (Command Prompt, Andrew Dunstan)

- In PL/Tcl, SPI commands are now run in subtransactions. If an error occurs, the subtransaction is cleaned up and the error is reported as an ordinary Tcl error, which can be trapped with catch. Formerly, it was not possible to catch such errors.

- Accept ELSEIF in PL/pgSQL (Neil)

  Previously PL/pgSQL only allowed ELSIF, but many people are accustomed to spelling this keyword ELSEIF.

**E.280.4.8. psql Changes**

- Improve psql information display about database objects (Christopher)

- Allow psql to display group membership in \du and \dg (Markus Bertheau)

- Prevent psql \dn from showing temporary schemas (Bruce)

- Allow psql to handle tilde user expansion for file names (Zach Irmen)

- Allow psql to display fancy prompts, including color, via readline (Reece Hart, Chet Ramey)

- Make psql \copy match COPY command syntax fully (Tom)

- Show the location of syntax errors (Fabien Coelho, Tom)

- Add CLUSTER information to psql \d display (Bruce)

- Change psql \copy stdin/stdout to read from command input/output (Bruce)

- Add pstdin/pstdout to read from psql's stdin/stdout (Mark Feit)

- Add global psql configuration file, psqlrc.sample (Bruce)

  This allows a central file where global psql startup commands can be stored.

- Have psql \d+ indicate if the table has an OID column (Neil)

- On Windows, use binary mode in psql when reading files so control-Z is not seen as end-of-file

- Have \dn+ show permissions and description for schemas (Dennis Björklund)

- Improve tab completion support (Stefan Kaltenbrunn, Greg Sabino Mullane)
• Allow boolean settings to be set using upper or lower case (Michael Paesold)

E.280.4.9. pg_dump Changes

• Use dependency information to improve the reliability of pg_dump (Tom)
  This should solve the longstanding problems with related objects sometimes being dumped in the
  wrong order.
• Have pg_dump output objects in alphabetical order if possible (Tom)
  This should make it easier to identify changes between dump files.
• Allow pg_restore to ignore some SQL errors (Fabien Coelho)
  This makes pg_restore's behavior similar to the results of feeding a pg_dump output script to psql.
  In most cases, ignoring errors and plowing ahead is the most useful thing to do. Also added was a
  pg_restore option to give the old behavior of exiting on an error.
• pg_restore --display now includes objects' schema names
• New begin/end markers in pg_dump text output (Bruce)
• Add start/stop times for pg_dump/pg_dumpall in verbose mode (Bruce)
• Allow most pg_dump options in pg_dumpall (Christopher)
• Have pg_dump use ALTER OWNER rather than SET SESSION AUTHORIZATION by default
  (Christopher)

E.280.4.10. libpq Changes

• Make libpq's SIGPIPE handling thread-safe (Bruce)
• Add PQmbdspalen() which returns the display length of a character (Tatsuo)
• Add thread locking to SSL and Kerberos connections (Manfred Spraul)
• Allow PQoidValue(), PQcmdTuples(), and PQoidStatus() to work on EXECUTE commands (Neil)
• Add PQserverVersion() to provide more convenient access to the server version number
  (Greg Sabino Mullane)
• Add PQprepare/PQsendPrepared() functions to support preparing statements without necessarily
  specifying the data types of their parameters (Abhijit Menon-Sen)
• Many ECPG improvements, including SET DESCRIPTOR (Michael)

E.280.4.11. Source Code Changes

• Allow the database server to run natively on Windows (Claudio, Magnus, Andrew)
• Shell script commands converted to C versions for Windows support (Andrew)
• Create an extension makefile framework (Fabien Coelho, Peter)
  This simplifies the task of building extensions outside the original source tree.
• Support relocatable installations (Bruce)
Directory paths for installed files (such as the /share directory) are now computed relative to the actual location of the executables, so that an installation tree can be moved to another place without reconfiguring and rebuilding.

- Use --with-docdir to choose installation location of documentation; also allow --infodir (Peter)
- Add --without-docdir to prevent installation of documentation (Peter)
- Upgrade to DocBook V4.2 SGML (Peter)
- New PostgreSQL CVS tag (Marc)

This was done to make it easier for organizations to manage their own copies of the PostgreSQL CVS repository. File version stamps from the master repository will not get munged by checking into or out of a copied repository.

- Clarify locking code (Manfred Koizar)
- Buffer manager cleanup (Neil)
- Decouple platform tests from CPU spinlock code (Bruce, Tom)
- Add inlined test-and-set code on PA-RISC for gcc (ViSolve, Tom)
- Improve i386 spinlock code (Manfred Spraul)
- Clean up spinlock assembly code to avoid warnings from newer gcc releases (Tom)
- Remove JDBC from source tree; now a separate project
- Remove the libpgtcl client interface; now a separate project
- More accurately estimate memory and file descriptor usage (Tom)
- Improvements to the macOS startup scripts (Ray A.)
- New fsync() test program (Bruce)
- Major documentation improvements (Neil, Peter)
- Remove pg_encoding; not needed anymore
- Remove pg_id; not needed anymore
- Remove initlocation; not needed anymore
- Auto-detect thread flags (no more manual testing) (Bruce)
- Use Olson's public domain timezone library (Magnus)
- With threading enabled, use thread flags on Unixware for backend executables too (Bruce)

Unixware cannot mix threaded and nonthreaded object files in the same executable, so everything must be compiled as threaded.

- psql now uses a flex-generated lexical analyzer to process command strings
- Reimplement the linked list data structure used throughout the backend (Neil)

This improves performance by allowing list append and length operations to be more efficient.

- Allow dynamically loaded modules to create their own server configuration parameters (Thomas Hallgren)
• New Brazilian version of FAQ (Euler Taveira de Oliveira)
• Add French FAQ (Guillaume Lelarge)
• New pgevent for Windows logging
• Make libpq and ECPG build as proper shared libraries on macOS (Tom)

E.280.4.12. Contrib Changes

• Overhaul of contrib/dblink (Joe)
• contrib/dbmirror improvements (Steven Singer)
• New contrib/xml2 (John Gray, Torchbox)
• Updated contrib/mysql
• New version of contrib/btree_gist (Teodor)
• New contrib/trgm, trigram matching for PostgreSQL (Teodor)
• Many contrib/tsearch2 improvements (Teodor)
• Add double metaphone to contrib/fuzzystrmatch (Andrew)
• Allow contrib/pg_autovacuum to run as a Windows service (Dave Page)
• Add functions to contrib/dbsize (Andreas Pflug)
• Removed contrib/pg_logger: obsoleted by integrated logging subprocess
• Removed contrib/rserv: obsoleted by various separate projects

E.281. Release 7.4.30

Release date: 2010-10-04

This release contains a variety of fixes from 7.4.29. For information about new features in the 7.4 major release, see Section E.311.

This is expected to be the last PostgreSQL release in the 7.4.X series. Users are encouraged to update to a newer release branch soon.

E.281.1. Migration to Version 7.4.30

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.26, see Section E.285.

E.281.2. Changes

• Use a separate interpreter for each calling SQL userid in PL/Perl and PL/Tcl (Tom Lane)

This change prevents security problems that can be caused by subverting Perl or Tcl code that will be executed later in the same session under another SQL user identity (for example, within a SECURITY DEFINER function). Most scripting languages offer numerous ways that that might be done, such as redefining standard functions or operators called by the target function. Without this change, any SQL user with Perl or Tcl language usage rights can do essentially anything with the SQL privileges of the target function’s owner.
The cost of this change is that intentional communication among Perl and Tcl functions becomes more difficult. To provide an escape hatch, PL/PerlU and PL/TclU functions continue to use only one interpreter per session. This is not considered a security issue since all such functions execute at the trust level of a database superuser already.

It is likely that third-party procedural languages that claim to offer trusted execution have similar security issues. We advise contacting the authors of any PL you are depending on for security-critical purposes.

Our thanks to Tim Bunce for pointing out this issue (CVE-2010-3433).

- Prevent possible crashes in `pg_get_expr()` by disallowing it from being called with an argument that is not one of the system catalog columns it's intended to be used with (Heikki Linnakangas, Tom Lane)

- Fix “cannot handle unplanned sub-select” error (Tom Lane)

  This occurred when a sub-select contains a join alias reference that expands into an expression containing another sub-select.

- Take care to fsync the contents of lockfiles (both `postmaster.pid` and the socket lockfile) while writing them (Tom Lane)

  This omission could result in corrupted lockfile contents if the machine crashes shortly after postmaster start. That could in turn prevent subsequent attempts to start the postmaster from succeeding, until the lockfile is manually removed.

- Improve `contrib/dblink`'s handling of tables containing dropped columns (Tom Lane)

- Fix connection leak after “duplicate connection name” errors in `contrib/dblink` (Itagaki Takahiro)

- Update build infrastructure and documentation to reflect the source code repository's move from CVS to Git (Magnus Hagander and others)

E.282. Release 7.4.29

Release date: 2010-05-17

This release contains a variety of fixes from 7.4.28. For information about new features in the 7.4 major release, see Section E.311.

The PostgreSQL community will stop releasing updates for the 7.4.X release series in July 2010. Users are encouraged to update to a newer release branch soon.

E.282.1. Migration to Version 7.4.29

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.26, see Section E.285.

E.282.2. Changes

- Enforce restrictions in `plperl` using an opmask applied to the whole interpreter, instead of using `Safe.pm` (Tim Bunce, Andrew Dunstan)

  Recent developments have convinced us that `Safe.pm` is too insecure to rely on for making `plperl` trustable. This change removes use of `Safe.pm` altogether, in favor of using a separate interpreter with an opcode mask that is always applied. Pleasant side effects of the change include
that it is now possible to use Perl's `strict` pragma in a natural way in `plperl`, and that Perl's `$a` and `$b` variables work as expected in sort routines, and that function compilation is significantly faster. (CVE-2010-1169)

- Prevent PL/Tcl from executing untrustworthy code from `pltcl_modules` (Tom)

  PL/Tcl's feature for autoloading Tcl code from a database table could be exploited for trojan-horse attacks, because there was no restriction on who could create or insert into that table. This change disables the feature unless `pltcl_modules` is owned by a superuser. (However, the permissions on the table are not checked, so installations that really need a less-than-secure modules table can still grant suitable privileges to trusted non-superusers.) Also, prevent loading code into the unrestricted “normal” Tcl interpreter unless we are really going to execute a `pltclu` function. (CVE-2010-1170)

- Do not allow an unprivileged user to reset superuser-only parameter settings (Alvaro)

  Previously, if an unprivileged user ran `ALTER USER ... RESET ALL` for himself, or `ALTER DATABASE ... RESET ALL` for a database he owns, this would remove all special parameter settings for the user or database, even ones that are only supposed to be changeable by a superuser. Now, the `ALTER` will only remove the parameters that the user has permission to change.

- Avoid possible crash during backend shutdown if shutdown occurs when a CONTEXT addition would be made to log entries (Tom)

  In some cases the context-printing function would fail because the current transaction had already been rolled back when it came time to print a log message.

- Update PL/Perl's `ppport.h` for modern Perl versions (Andrew)

- Fix assorted memory leaks in PL/Python (Andreas Freund, Tom)

- Ensure that `contrib/pgstattuple` functions respond to cancel interrupts promptly (Tatsuhito Kasahara)

- Make server startup deal properly with the case that `shmget()` returns `EINVAL` for an existing shared memory segment (Tom)

  This behavior has been observed on BSD-derived kernels including macOS. It resulted in an entirely-misleading startup failure complaining that the shared memory request size was too large.

**E.283. Release 7.4.28**

Release date: 2010-03-15

This release contains a variety of fixes from 7.4.27. For information about new features in the 7.4 major release, see Section E.311.

The PostgreSQL community will stop releasing updates for the 7.4.X release series in July 2010. Users are encouraged to update to a newer release branch soon.

**E.283.1. Migration to Version 7.4.28**

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.26, see Section E.285.

**E.283.2. Changes**

- Add new configuration parameter `ssl_renegotiation_limit` to control how often we do session key renegotiation for an SSL connection (Magnus)
This can be set to zero to disable renegotiation completely, which may be required if a broken SSL library is used. In particular, some vendors are shipping stopgap patches for CVE-2009-3555 that cause renegotiation attempts to fail.

- Make substring() for bit types treat any negative length as meaning “all the rest of the string” (Tom)

  The previous coding treated only -1 that way, and would produce an invalid result value for other negative values, possibly leading to a crash (CVE-2010-0442).

- Fix some cases of pathologically slow regular expression matching (Tom)

- When reading pg_hba.conf and related files, do not treat @something as a file inclusion request if the @ appears inside quote marks; also, never treat @ by itself as a file inclusion request (Tom)

  This prevents erratic behavior if a role or database name starts with @. If you need to include a file whose path name contains spaces, you can still do so, but you must write @"/path to/file" rather than putting the quotes around the whole construct.

- Prevent infinite loop on some platforms if a directory is named as an inclusion target in pg_hba.conf and related files (Tom)

- Ensure PL/Tcl initializes the Tcl interpreter fully (Tom)

  The only known symptom of this oversight is that the Tcl clock command misbehaves if using Tcl 8.5 or later.

- Prevent crash in contrib/dblink when too many key columns are specified to a dblink_build_sql_* function (Rushabh Lathia, Joe Conway)

**E.284. Release 7.4.27**

**Release date:** 2009-12-14

This release contains a variety of fixes from 7.4.26. For information about new features in the 7.4 major release, see Section E.311.

**E.284.1. Migration to Version 7.4.27**

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.26, see Section E.285.

**E.284.2. Changes**

- Protect against indirect security threats caused by index functions changing session-local state (Gurjeet Singh, Tom)

  This change prevents allegedly-immutable index functions from possibly subverting a superuser's session (CVE-2009-4136).

- Reject SSL certificates containing an embedded null byte in the common name (CN) field (Magnus)

  This prevents unintended matching of a certificate to a server or client name during SSL validation (CVE-2009-4034).

- Fix possible crash during backend-startup-time cache initialization (Tom)
• Prevent signals from interrupting VACUUM at unsafe times (Alvaro)

This fix prevents a PANIC if a VACUUM FULL is canceled after it's already committed its tuple movements, as well as transient errors if a plain VACUUM is interrupted after having truncated the table.

• Fix possible crash due to integer overflow in hash table size calculation (Tom)

This could occur with extremely large planner estimates for the size of a hashjoin’s result.

• Fix very rare crash in inet/cidr comparisons (Chris Mikkelson)

• Fix PAM password processing to be more robust (Tom)

The previous code is known to fail with the combination of the Linux pam_krb5 PAM module with Microsoft Active Directory as the domain controller. It might have problems elsewhere too, since it was making unjustified assumptions about what arguments the PAM stack would pass to it.

• Make the postmaster ignore any application_name parameter in connection request packets, to improve compatibility with future libpq versions (Tom)

E.285. Release 7.4.26

Release date: 2009-09-09

This release contains a variety of fixes from 7.4.25. For information about new features in the 7.4 major release, see Section E.311.

E.285.1. Migration to Version 7.4.26

A dump/restore is not required for those running 7.4.X. However, if you have any hash indexes on interval columns, you must REINDEX them after updating to 7.4.26. Also, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.285.2. Changes

• Disallow RESET ROLE and RESET SESSION AUTHORIZATION inside security-definer functions (Tom, Heikki)

This covers a case that was missed in the previous patch that disallowed SET ROLE and SET SESSION AUTHORIZATION inside security-definer functions. (See CVE-2007-6600)

• Fix handling of sub-SELECTs appearing in the arguments of an outer-level aggregate function (Tom)

• Fix hash calculation for data type interval (Tom)

This corrects wrong results for hash joins on interval values. It also changes the contents of hash indexes on interval columns. If you have any such indexes, you must REINDEX them after updating.

• Fix overflow for INTERVAL 'x ms' when x is more than 2 million and integer datetimes are in use (Alex Hunsaker)

• Fix calculation of distance between a point and a line segment (Tom)

This led to incorrect results from a number of geometric operators.

• Fix money data type to work in locales where currency amounts have no fractional digits, e.g. Japan (Itagaki Takahiro)
• Properly round datetime input like 00:12:57.9999999999999999999999999999 (Tom)
• Fix poor choice of page split point in GiST R-tree operator classes (Teodor)
• Fix portability issues in plperl initialization (Andrew Dunstan)
• Improve robustness of libpq's code to recover from errors during COPY FROM STDIN (Tom)
• Avoid including conflicting readline and editline header files when both libraries are installed (Zdenek Kotala)

E.286. Release 7.4.25

Release date: 2009-03-16

This release contains a variety of fixes from 7.4.24. For information about new features in the 7.4 major release, see Section E.311.

E.286.1. Migration to Version 7.4.25

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.286.2. Changes

• Prevent error recursion crashes when encoding conversion fails (Tom)

  This change extends fixes made in the last two minor releases for related failure scenarios. The previous fixes were narrowly tailored for the original problem reports, but we have now recognized that any error thrown by an encoding conversion function could potentially lead to infinite recursion while trying to report the error. The solution therefore is to disable translation and encoding conversion and report the plain-ASCII form of any error message, if we find we have gotten into a recursive error reporting situation. (CVE-2009-0922)

• Disallow CREATE CONVERSION with the wrong encodings for the specified conversion function (Heikki)

  This prevents one possible scenario for encoding conversion failure. The previous change is a backstop to guard against other kinds of failures in the same area.

• Fix core dump when to_char() is given format codes that are inappropriate for the type of the data argument (Tom)

• Add MUST (Mauritius Island Summer Time) to the default list of known timezone abbreviations (Xavier Bugaud)

E.287. Release 7.4.24

Release date: 2009-02-02

This release contains a variety of fixes from 7.4.23. For information about new features in the 7.4 major release, see Section E.311.

E.287.1. Migration to Version 7.4.24

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.
E.287.2. Changes

- Improve handling of URLs in `headline()` function (Teodor)
- Improve handling of overlength headlines in `headline()` function (Teodor)
- Prevent possible Assert failure or misconversion if an encoding conversion is created with the wrong conversion function for the specified pair of encodings (Tom, Heikki)
- Avoid unnecessary locking of small tables in `VACUUM` (Heikki)
- Fix uninitialized variables in `contrib/tsearch2`'s `get_covers()` function (Teodor)
- Fix bug in `to_char`'s handling of TH format codes (Andreas Scherbaum)
- Make all documentation reference `pgsql-bugs` and/or `pgsql-hackers` as appropriate, instead of the now-decommissioned `pgsql-ports` and `pgsql-patches` mailing lists (Tom)

E.288. Release 7.4.23

**Release date:** 2008-11-03

This release contains a variety of fixes from 7.4.22. For information about new features in the 7.4 major release, see Section E.311.

E.288.1. Migration to Version 7.4.23

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.288.2. Changes

- Fix backend crash when the client encoding cannot represent a localized error message (Tom)
  
  We have addressed similar issues before, but it would still fail if the “character has no equivalent” message itself couldn’t be converted. The fix is to disable localization and send the plain ASCII error message when we detect such a situation.

- Fix incorrect tssearch2 headline generation when single query item matches first word of text (Sushant Sinha)

- Fix improper display of fractional seconds in interval values when using a non-ISO datestyle in an `--enable-integer-datetimes` build (Ron Mayer)

- Ensure `SPI_getvalue` and `SPI_getbinval` behave correctly when the passed tuple and tuple descriptor have different numbers of columns (Tom)
  
  This situation is normal when a table has had columns added or removed, but these two functions didn’t handle it properly. The only likely consequence is an incorrect error indication.

- Fix ecpg's parsing of `CREATE USER` (Michael)

E.289. Release 7.4.22

**Release date:** 2008-09-22

This release contains a variety of fixes from 7.4.21. For information about new features in the 7.4 major release, see Section E.311.
E.289.1. Migration to Version 7.4.22

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.289.2. Changes

- Fix datetime input functions to correctly detect integer overflow when running on a 64-bit platform (Tom)
- Improve performance of writing very long log messages to syslog (Tom)
- Fix bug in backwards scanning of a cursor on a SELECT DISTINCT ON query (Tom)
- Fix planner to estimate that GROUP BY expressions yielding boolean results always result in two groups, regardless of the expressions' contents (Tom)

This is very substantially more accurate than the regular GROUP BY estimate for certain boolean tests like col IS NULL.

- Improve pg_dump and pg_restore's error reporting after failure to send a SQL command (Tom)

E.290. Release 7.4.21

Release date: 2008-06-12

This release contains one serious bug fix over 7.4.20. For information about new features in the 7.4 major release, see Section E.311.

E.290.1. Migration to Version 7.4.21

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.290.2. Changes

- Make pg_get_ruledef() parenthesize negative constants (Tom)

Before this fix, a negative constant in a view or rule might be dumped as, say, -42::integer, which is subtly incorrect: it should be (-42)::integer due to operator precedence rules. Usually this would make little difference, but it could interact with another recent patch to cause PostgreSQL to reject what had been a valid SELECT DISTINCT view query. Since this could result in pg_dump output failing to reload, it is being treated as a high-priority fix. The only released versions in which dump output is actually incorrect are 8.3.1 and 8.2.7.

E.291. Release 7.4.20

Release date: never released

This release contains a variety of fixes from 7.4.19. For information about new features in the 7.4 major release, see Section E.311.

E.291.1. Migration to Version 7.4.20

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.
E.291.2. Changes

- Fix conversions between ISO-8859-5 and other encodings to handle Cyrillic “Yo” characters (ё and Е with two dots) (Sergey Burladyan)

- Fix a few datatype input functions that were allowing unused bytes in their results to contain uninitialized, unpredictable values (Tom)

  This could lead to failures in which two apparently identical literal values were not seen as equal, resulting in the parser complaining about unmatched ORDER BY and DISTINCT expressions.

- Fix a corner case in regular-expression substring matching (substring(string from pattern)) (Tom)

  The problem occurs when there is a match to the pattern overall but the user has specified a parenthesized subexpression and that subexpression hasn't got a match. An example is substring('foo' from 'foo(bar)?'). This should return NULL, since (bar) isn't matched, but it was mistakenly returning the whole-pattern match instead (ie, foo).

- Fix incorrect result from ecpg's PGTYPEStimestamp_sub() function (Michael)

- Fix DatumGetBool macro to not fail with gcc 4.3 (Tom)

  This problem affects “old style” (V0) C functions that return boolean. The fix is already in 8.3, but the need to back-patch it was not realized at the time.

- Fix longstanding LISTEN/NOTIFY race condition (Tom)

  In rare cases a session that had just executed a LISTEN might not get a notification, even though one would be expected because the concurrent transaction executing NOTIFY was observed to commit later.

  A side effect of the fix is that a transaction that has executed a not-yet-committed LISTEN command will not see any row in pg_listener for the LISTEN, should it choose to look; formerly it would have. This behavior was never documented one way or the other, but it is possible that some applications depend on the old behavior.

- Fix display of constant expressions in ORDER BY and GROUP BY (Tom)

  An explicitly casted constant would be shown incorrectly. This could for example lead to corruption of a view definition during dump and reload.

- Fix libpq to handle NOTICE messages correctly during COPY OUT (Tom)

  This failure has only been observed to occur when a user-defined datatype's output routine issues a NOTICE, but there is no guarantee it couldn't happen due to other causes.

E.292. Release 7.4.19

Release date: 2008-01-07

This release contains a variety of fixes from 7.4.18, including fixes for significant security issues. For information about new features in the 7.4 major release, see Section E.311.

E.292.1. Migration to Version 7.4.19

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.
E.292.2. Changes

• Prevent functions in indexes from executing with the privileges of the user running VACUUM, ANALYZE, etc (Tom)

Functions used in index expressions and partial-index predicates are evaluated whenever a new table entry is made. It has long been understood that this poses a risk of trojan-horse code execution if one modifies a table owned by an untrustworthy user. (Note that triggers, defaults, check constraints, etc. pose the same type of risk.) But functions in indexes pose extra danger because they will be executed by routine maintenance operations such as VACUUM FULL, which are commonly performed automatically under a superuser account. For example, a nefarious user can execute code with superuser privileges by setting up a trojan-horse index definition and waiting for the next routine vacuum. The fix arranges for standard maintenance operations (including VACUUM, ANALYZE, REINDEX, and CLUSTER) to execute as the table owner rather than the calling user, using the same privilege-switching mechanism already used for SECURITY DEFINER functions. To prevent bypassing this security measure, execution of SET SESSION AUTHORIZATION and SET ROLE is now forbidden within a SECURITY DEFINER context. (CVE-2007-6600)

• Repair assorted bugs in the regular-expression package (Tom, Will Drewry)

Suitably crafted regular-expression patterns could cause crashes, infinite or near-infinite looping, and/or massive memory consumption, all of which pose denial-of-service hazards for applications that accept regex search patterns from untrustworthy sources. (CVE-2007-4769, CVE-2007-4772, CVE-2007-6067)

• Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

The fix that appeared for this in 7.4.18 was incomplete, as it plugged the hole for only some dblink functions. (CVE-2007-6601, CVE-2007-3278)

• Fix planner failure in some cases of WHERE false AND var IN (SELECT ...) (Tom)

• Fix potential crash in translate() when using a multibyte database encoding (Tom)

• Fix PL/Python to not crash on long exception messages (Alvaro)

• ecpg parser fixes (Michael)

• Make contrib/tablefunc's crosstab() handle NULL rowid as a category in its own right, rather than crashing (Joe)

• Fix tsvector and tsquery output routines to escape backslashes correctly (Teodor, Bruce)

• Fix crash of to_tsvector() on huge input strings (Teodor)

• Require a specific version of Autoconf to be used when re-generating the configure script (Peter)

This affects developers and packagers only. The change was made to prevent accidental use of untested combinations of Autoconf and PostgreSQL versions. You can remove the version check if you really want to use a different Autoconf version, but it's your responsibility whether the result works or not.

E.293. Release 7.4.18

Release date: 2007-09-17

This release contains fixes from 7.4.17. For information about new features in the 7.4 major release, see Section E.311.
E.293.1. Migration to Version 7.4.18

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.293.2. Changes

- Prevent index corruption when a transaction inserts rows and then aborts close to the end of a concurrent VACUUM on the same table (Tom)
- Make CREATE DOMAIN ... DEFAULT NULL work properly (Tom)
- Fix excessive logging of SSL error messages (Tom)
- Fix crash when log_min_error_statement logging runs out of memory (Tom)
- Prevent CLUSTER from failing due to attempting to process temporary tables of other sessions (Alvaro)
- Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

E.294. Release 7.4.17

Release date: 2007-04-23

This release contains fixes from 7.4.16, including a security fix. For information about new features in the 7.4 major release, see Section E.311.

E.294.1. Migration to Version 7.4.17

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.294.2. Changes

- Support explicit placement of the temporary-table schema within search_path, and disable searching it for functions and operators (Tom)
  This is needed to allow a security-definer function to set a truly secure value of search_path. Without it, an unprivileged SQL user can use temporary objects to execute code with the privileges of the security-definer function (CVE-2007-2138). See CREATE FUNCTION for more information.
- /contrib/tsearch2 crash fixes (Teodor)
- Fix potential-data-corruption bug in how VACUUM FULL handles UPDATE chains (Tom, Pavan Deolasee)
- Fix PANIC during enlargement of a hash index (bug introduced in 7.4.15) (Tom)

E.295. Release 7.4.16

Release date: 2007-02-05

This release contains a variety of fixes from 7.4.15, including a security fix. For information about new features in the 7.4 major release, see Section E.311.
E.295.1. Migration to Version 7.4.16

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.295.2. Changes

• Remove security vulnerability that allowed connected users to read backend memory (Tom)

  The vulnerability involves suppressing the normal check that a SQL function returns the data type it's declared to, or changing the data type of a table column used in a SQL function (CVE-2007-0555). This error can easily be exploited to cause a backend crash, and in principle might be used to read database content that the user should not be able to access.

• Fix rare bug wherein btree index page splits could fail due to choosing an infeasible split point (Heikki Linnakangas)

• Fix for rare Assert() crash triggered by UNION (Tom)

• Tighten security of multi-byte character processing for UTF8 sequences over three bytes long (Tom)

E.296. Release 7.4.15

Release date: 2007-01-08

This release contains a variety of fixes from 7.4.14. For information about new features in the 7.4 major release, see Section E.311.

E.296.1. Migration to Version 7.4.15

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.296.2. Changes

• Improve handling of getaddrinfo() on AIX (Tom)

  This fixes a problem with starting the statistics collector, among other things.

• Fix “failed to re-find parent key” errors in VACUUM (Tom)

• Fix bugs affecting multi-gigabyte hash indexes (Tom)

• Fix error when constructing an ARRAY[] made up of multiple empty elements (Tom)

• to_number() and to_char(numeric) are now STABLE, not IMMUTABLE, for new initdb installs (Tom)

  This is because lc_numeric can potentially change the output of these functions.

• Improve index usage of regular expressions that use parentheses (Tom)

  This improves psql \d performance also.

E.297. Release 7.4.14

Release date: 2006-10-16
This release contains a variety of fixes from 7.4.13. For information about new features in the 7.4 major release, see Section E.311.

E.297.1. Migration to Version 7.4.14

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

E.297.2. Changes

- Fix core dump when an untyped literal is taken as ANYARRAY
- Fix `string_to_array()` to handle overlapping matches for the separator string
  For example, `string_to_array('123xx456xxx789', 'xx').`
- Fix corner cases in pattern matching for psql's `\d` commands
- Fix index-corrupting bugs in `/contrib/ltree` (Teodor)
- Fix backslash escaping in `/contrib/dbmirror`
- Adjust regression tests for recent changes in US DST laws

E.298. Release 7.4.13

Release date: 2006-05-23

This release contains a variety of fixes from 7.4.12, including patches for extremely serious security issues. For information about new features in the 7.4 major release, see Section E.311.

E.298.1. Migration to Version 7.4.13

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

Full security against the SQL-injection attacks described in CVE-2006-2313 and CVE-2006-2314 might require changes in application code. If you have applications that embed untrustworthy strings into SQL commands, you should examine them as soon as possible to ensure that they are using recommended escaping techniques. In most cases, applications should be using subroutines provided by libraries or drivers (such as libpq's `PQescapeStringConn()`) to perform string escaping, rather than relying on ad hoc code to do it.

E.298.2. Changes

- Change the server to reject invalidly-encoded multibyte characters in all cases (Tatsuo, Tom)
  While PostgreSQL has been moving in this direction for some time, the checks are now applied uniformly to all encodings and all textual input, and are now always errors not merely warnings. This change defends against SQL-injection attacks of the type described in CVE-2006-2313.

- Reject unsafe uses of `\'` in string literals
  As a server-side defense against SQL-injection attacks of the type described in CVE-2006-2314, the server now only accepts `''` and not `\'` as a representation of ASCII single quote in SQL string literals. By default, `\'` is rejected only when `client_encoding` is set to a client-only encoding (SJIS, BIG5, GBK, GB18030, or UHC), which is the scenario in which SQL injection is possible. A
new configuration parameter `backslash_quote` is available to adjust this behavior when needed. Note that full security against CVE-2006-2314 might require client-side changes; the purpose of `backslash_quote` is in part to make it obvious that insecure clients are insecure.

- Modify libpq's string-escaping routines to be aware of encoding considerations and standard_conforming_strings

This fixes libpq-using applications for the security issues described in CVE-2006-2313 and CVE-2006-2314, and also future-proofs them against the planned changeover to SQL-standard string literal syntax. Applications that use multiple PostgreSQL connections concurrently should migrate to `PQescapeStringConn()` and `PQescapeByteaConn()` to ensure that escaping is done correctly for the settings in use in each database connection. Applications that do string escaping “by hand” should be modified to rely on library routines instead.

- Fix some incorrect encoding conversion functions
  
  `win1251_to_iso`, `alt_to_iso`, `euc_tw_to_big5`, `euc_tw_to_mic`, `mic_to_euc_tw` were all broken to varying extents.

- Clean up stray remaining uses of `\'` in strings (Bruce, Jan)

- Fix bug that sometimes caused OR’d index scans to miss rows they should have returned

- Fix WAL replay for case where a btree index has been truncated

- Fix `SIMILAR TO` for patterns involving `|` (Tom)

- Fix server to use custom DH SSL parameters correctly (Michael Fuhr)

- Fix for Bonjour on Intel Macs (Ashley Clark)

- Fix various minor memory leaks

**E.299. Release 7.4.12**

**Release date:** 2006-02-14

This release contains a variety of fixes from 7.4.11. For information about new features in the 7.4 major release, see Section E.311.

**E.299.1. Migration to Version 7.4.12**

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.11, see Section E.300.

**E.299.2. Changes**

- Fix potential crash in `SET SESSION AUTHORIZATION` (CVE-2006-0553)

  An unprivileged user could crash the server process, resulting in momentary denial of service to other users, if the server has been compiled with Asserts enabled (which is not the default). Thanks to Akio Ishida for reporting this problem.

- Fix bug with row visibility logic in self-inserted rows (Tom)

  Under rare circumstances a row inserted by the current command could be seen as already valid, when it should not be. Repairs bug created in 7.4.9 and 7.3.11 releases.

- Fix race condition that could lead to “file already exists” errors during `pg_clog` file creation (Tom)
• Properly check `DOMAIN` constraints for `UNKNOWN` parameters in prepared statements (Neil)
• Fix to allow restoring dumps that have cross-schema references to custom operators (Tom)
• Portability fix for testing presence of `finite` and `isinf` during configure (Tom)

E.300. Release 7.4.11

Release date: 2006-01-09

This release contains a variety of fixes from 7.4.10. For information about new features in the 7.4 major release, see Section E.311.

E.300.1. Migration to Version 7.4.11

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.8, see Section E.303. Also, you might need to `REINDEX` indexes on textual columns after updating, if you are affected by the locale or plperl issues described below.

E.300.2. Changes

• Fix for protocol-level Describe messages issued outside a transaction or in a failed transaction (Tom)
• Fix character string comparison for locales that consider different character combinations as equal, such as Hungarian (Tom)

This might require `REINDEX` to fix existing indexes on textual columns.
• Set locale environment variables during postmaster startup to ensure that plperl won't change the locale later

This fixes a problem that occurred if the postmaster was started with environment variables specifying a different locale than what initdb had been told. Under these conditions, any use of plperl was likely to lead to corrupt indexes. You might need `REINDEX` to fix existing indexes on textual columns if this has happened to you.
• Fix longstanding bug in `strpos()` and regular expression handling in certain rarely used Asian multi-byte character sets (Tatsuo)
• Fix bug in `/contrib/pgcrypto gen_salt`, which caused it not to use all available salt space for MD5 and XDES algorithms (Marko Kreen, Solar Designer)

Salts for Blowfish and standard DES are unaffected.
• Fix `/contrib/dblink` to throw an error, rather than crashing, when the number of columns specified is different from what's actually returned by the query (Joe)

E.301. Release 7.4.10

Release date: 2005-12-12

This release contains a variety of fixes from 7.4.9. For information about new features in the 7.4 major release, see Section E.311.

E.301.1. Migration to Version 7.4.10

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.8, see Section E.303.
E.301.2. Changes

- Fix race condition in transaction log management

  There was a narrow window in which an I/O operation could be initiated for the wrong page, leading to an Assert failure or data corruption.

- Prevent failure if client sends Bind protocol message when current transaction is already aborted

- /contrib/ltree fixes (Teodor)

- AIX and HPUX compile fixes (Tom)

- Fix longstanding planning error for outer joins

  This bug sometimes caused a bogus error “RIGHT JOIN is only supported with merge-joinable join conditions”.

- Prevent core dump in pg_autovacuum when a table has been dropped

E.302. Release 7.4.9

Release date: 2005-10-04

This release contains a variety of fixes from 7.4.8. For information about new features in the 7.4 major release, see Section E.311.

E.302.1. Migration to Version 7.4.9

A dump/restore is not required for those running 7.4.X. However, if you are upgrading from a version earlier than 7.4.8, see Section E.303.

E.302.2. Changes

- Fix error that allowed VACUUM to remove ctid chains too soon, and add more checking in code that follows ctid links

  This fixes a long-standing problem that could cause crashes in very rare circumstances.

- Fix CHAR() to properly pad spaces to the specified length when using a multiple-byte character set (Yoshiyuki Asaba)

  In prior releases, the padding of CHAR() was incorrect because it only padded to the specified number of bytes without considering how many characters were stored.

- Fix the sense of the test for read-only transaction in COPY

  The code formerly prohibited COPY TO, where it should prohibit COPY FROM.

- Fix planning problem with outer-join ON clauses that reference only the inner-side relation

- Further fixes for x FULL JOIN y ON true corner cases

- Make array_in and array_recv more paranoid about validating their OID parameter

- Fix missing rows in queries like UPDATE a=... WHERE a... with GiST index on column a

- Improve robustness of datetime parsing

- Improve checking for partially-written WAL pages
• Improve robustness of signal handling when SSL is enabled
• Don’t try to open more than max_files_per_process files during postmaster startup
• Various memory leakage fixes
• Various portability improvements
• Fix PL/pgSQL to handle var := var correctly when the variable is of pass-by-reference type
• Update contrib/tsearch2 to use current Snowball code

E.303. Release 7.4.8

Release date: 2005-05-09

This release contains a variety of fixes from 7.4.7, including several security-related issues. For information about new features in the 7.4 major release, see Section E.311.

E.303.1. Migration to Version 7.4.8

A dump/restore is not required for those running 7.4.X. However, it is one possible way of handling two significant security problems that have been found in the initial contents of 7.4.X system catalogs. A dump/initdb/reload sequence using 7.4.8’s initdb will automatically correct these problems.

The larger security problem is that the built-in character set encoding conversion functions can be invoked from SQL commands by unprivileged users, but the functions were not designed for such use and are not secure against malicious choices of arguments. The fix involves changing the declared parameter list of these functions so that they can no longer be invoked from SQL commands. (This does not affect their normal use by the encoding conversion machinery.)

The lesser problem is that the contrib/tsearch2 module creates several functions that are mis-declared to return internal when they do not accept internal arguments. This breaks type safety for all functions using internal arguments.

It is strongly recommended that all installations repair these errors, either by initdb or by following the manual repair procedures given below. The errors at least allow unprivileged database users to crash their server process, and might allow unprivileged users to gain the privileges of a database superuser.

If you wish not to do an initdb, perform the following procedures instead. As the database superuser, do:

BEGIN;
UPDATE pg_proc SET proargtypes[3] = 'internal'::regtype
WHERE pronamespace = 11 AND pronargs = 5
    AND proargtypes[2] = 'cstring'::regtype;
-- The command should report having updated 90 rows;
-- if not, rollback and investigate instead of committing!
COMMIT;

Next, if you have installed contrib/tsearch2, do:

BEGIN;
UPDATE pg_proc SET proargtypes[0] = 'internal'::regtype
WHERE oid IN ('dex_init(text) '::regprocedure,
              'snb_en_init(text) '::regprocedure,
'snb_ru_init(text)'::regprocedure,
'spell_init(text)'::regprocedure,
'syn_init(text)'::regprocedure
);
-- The command should report having updated 5 rows;
-- if not, rollback and investigate instead of committing!
COMMIT;

If this command fails with a message like “function "dex_init(text)" does not exist”, then either
tsearch2 is not installed in this database, or you already did the update.

The above procedures must be carried out in each database of an installation, including template1,
and ideally including template0 as well. If you do not fix the template databases then any subse-
quently created databases will contain the same errors. template1 can be fixed in the same way as
any other database, but fixing template0 requires additional steps. First, from any database issue:

UPDATE pg_database SET datallowconn = true WHERE datname =
'template0';

Next connect to template0 and perform the above repair procedures. Finally, do:

-- re-freeze template0:
VACUUM FREEZE;
-- and protect it against future alterations:
UPDATE pg_database SET datallowconn = false WHERE datname =
'template0';

E.303.2. Changes

• Change encoding function signature to prevent misuse

• Change contrib/tsearch2 to avoid unsafe use of INTERNAL function results

• Repair ancient race condition that allowed a transaction to be seen as committed for some purposes
  (eg SELECT FOR UPDATE) slightly sooner than for other purposes

  This is an extremely serious bug since it could lead to apparent data inconsistencies being briefly
  visible to applications.

• Repair race condition between relation extension and VACUUM

  This could theoretically have caused loss of a page's worth of freshly-inserted data, although the
  scenario seems of very low probability. There are no known cases of it having caused more than
  an Assert failure.

• Fix comparisons of TIME WITH TIME ZONE values

  The comparison code was wrong in the case where the --enable-integer-datetimes configuration switch had been used. NOTE: if you have an index on a TIME WITH TIME ZONE column, it will need to be REINDEXed after installing this update, because the fix corrects the sort order of column values.

• Fix EXTRACT(EPOCH) for TIME WITH TIME ZONE values

• Fix mis-display of negative fractional seconds in INTERVAL values

  This error only occurred when the --enable-integer-datetimes configuration switch had been used.
• Ensure operations done during backend shutdown are counted by statistics collector

This is expected to resolve reports of pg_autovacuum not vacuuming the system catalogs often enough — it was not being told about catalog deletions caused by temporary table removal during backend exit.

• Additional buffer overrun checks in plpgsql (Neil)

• Fix pg_dump to dump trigger names containing % correctly (Neil)

• Fix contrib/pgcrypto for newer OpenSSL builds (Marko Kreen)

• Still more 64-bit fixes for contrib/intagg

• Prevent incorrect optimization of functions returning RECORD

• Prevent to_char(interval) from dumping core for month-related formats

• Prevent crash on COALESCE(NULL, NULL)

• Fix array_map to call PL functions correctly

• Fix permission checking in ALTER DATABASE RENAME

• Fix ALTER LANGUAGE RENAME

• Make RemoveFromWaitQueue clean up after itself

This fixes a lock management error that would only be visible if a transaction was kicked out of a wait for a lock (typically by query cancel) and then the holder of the lock released it within a very narrow window.

• Fix problem with untyped parameter appearing in INSERT ... SELECT

• Fix CLUSTER failure after ALTER TABLE SET WITHOUT OIDS

E.304. Release 7.4.7

Release date: 2005-01-31

This release contains a variety of fixes from 7.4.6, including several security-related issues. For information about new features in the 7.4 major release, see Section E.311.

E.304.1. Migration to Version 7.4.7

A dump/restore is not required for those running 7.4.X.

E.304.2. Changes

• Disallow LOAD to non-superusers

On platforms that will automatically execute initialization functions of a shared library (this includes at least Windows and ELF-based Unixen), LOAD can be used to make the server execute arbitrary code. Thanks to NGS Software for reporting this.

• Check that creator of an aggregate function has the right to execute the specified transition functions

This oversight made it possible to bypass denial of EXECUTE permission on a function.

• Fix security and 64-bit issues in contrib/intagg
• Add needed STRICT marking to some contrib functions (Kris Jurka)
• Avoid buffer overrun when plpgsql cursor declaration has too many parameters (Neil)
• Fix planning error for FULL and RIGHT outer joins
  The result of the join was mistakenly supposed to be sorted the same as the left input. This could
  not only deliver mis-sorted output to the user, but in case of nested merge joins could give outright
  wrong answers.
• Fix plperl for quote marks in tuple fields
• Fix display of negative intervals in SQL and GERMAN datestyles
• Make age(timestamptz) do calculation in local timezone not GMT

E.305. Release 7.4.6

Release date: 2004-10-22

This release contains a variety of fixes from 7.4.5. For information about new features in the 7.4 major
release, see Section E.311.

E.305.1. Migration to Version 7.4.6

A dump/restore is not required for those running 7.4.X.

E.305.2. Changes

• Repair possible failure to update hint bits on disk
  Under rare circumstances this oversight could lead to “could not access transaction status” failures,
  which qualifies it as a potential-data-loss bug.
• Ensure that hashed outer join does not miss tuples
  Very large left joins using a hash join plan could fail to output unmatched left-side rows given just
  the right data distribution.
• Disallow running pg_ctl as root
  This is to guard against any possible security issues.
• Avoid using temp files in /tmp in make_oidjoins_check
  This has been reported as a security issue, though it's hardly worthy of concern since there is no
  reason for non-developers to use this script anyway.
• Prevent forced backend shutdown from re-emitting prior command result
  In rare cases, a client might think that its last command had succeeded when it really had been
  aborted by forced database shutdown.
• Repair bug in pg_stat_get_backend_idset
  This could lead to misbehavior in some of the system-statistics views.
• Fix small memory leak in postmaster
• Fix “expected both swapped tables to have TOAST tables” bug
This could arise in cases such as CLUSTER after ALTER TABLE DROP COLUMN.

- Prevent `pg_ctl restart` from adding `-D` multiple times
- Fix problem with NULL values in GiST indexes
- `::` is no longer interpreted as a variable in an ECPG prepare statement

**E.306. Release 7.4.5**

**Release date:** 2004-08-18

This release contains one serious bug fix over 7.4.4. For information about new features in the 7.4 major release, see Section E.311.

**E.306.1. Migration to Version 7.4.5**

A dump/restore is not required for those running 7.4.X.

**E.306.2. Changes**

- Repair possible crash during concurrent B-tree index insertions

  This patch fixes a rare case in which concurrent insertions into a B-tree index could result in a server panic. No permanent damage would result, but it's still worth a re-release. The bug does not exist in pre-7.4 releases.

**E.307. Release 7.4.4**

**Release date:** 2004-08-16

This release contains a variety of fixes from 7.4.3. For information about new features in the 7.4 major release, see Section E.311.

**E.307.1. Migration to Version 7.4.4**

A dump/restore is not required for those running 7.4.X.

**E.307.2. Changes**

- Prevent possible loss of committed transactions during crash

  Due to insufficient interlocking between transaction commit and checkpointing, it was possible for transactions committed just before the most recent checkpoint to be lost, in whole or in part, following a database crash and restart. This is a serious bug that has existed since PostgreSQL 7.1.

- Check HAVING restriction before evaluating result list of an aggregate plan

- Avoid crash when session's current user ID is deleted

- Fix hashed crosstab for zero-rows case (Joe)

- Force cache update after renaming a column in a foreign key

- Pretty-print UNION queries correctly

- Make `psql` handle `\r\n` newlines properly in COPY IN
- pg_dump handled ACLs with grant options incorrectly
- Fix thread support for macOS and Solaris
- Updated JDBC driver (build 215) with various fixes
- ECPG fixes
- Translation updates (various contributors)

**E.308. Release 7.4.3**

**Release date:** 2004-06-14

This release contains a variety of fixes from 7.4.2. For information about new features in the 7.4 major release, see Section E.311.

**E.308.1. Migration to Version 7.4.3**

A dump/restore is not required for those running 7.4.X.

**E.308.2. Changes**

- Fix temporary memory leak when using non-hashed aggregates (Tom)
- ECPG fixes, including some for Informix compatibility (Michael)
- Fixes for compiling with thread-safety, particularly Solaris (Bruce)
- Fix error in COPY IN termination when using the old network protocol (ljb)
- Several important fixes in pg_autovacuum, including fixes for large tables, unsigned oids, stability, temp tables, and debug mode (Matthew T. O’Connor)
- Fix problem with reading tar-format dumps on NetBSD and BSD/OS (Bruce)
- Several JDBC fixes
- Fix ALTER SEQUENCE RESTART where last_value equals the restart value (Tom)
- Repair failure to recalculate nested sub-selects (Tom)
- Fix problems with non-constant expressions in LIMIT/OFFSET
- Support FULL JOIN with no join clause, such as X FULL JOIN Y ON TRUE (Tom)
- Fix another zero-column table bug (Tom)
- Improve handling of non-qualified identifiers in GROUP BY clauses in sub-selects (Tom)
- Select-list aliases within the sub-select will now take precedence over names from outer query levels.
- Do not generate “NATURAL CROSS JOIN” when decompiling rules (Tom)
- Add checks for invalid field length in binary COPY (Tom)
- This fixes a difficult-to-exploit security hole.
- Avoid locking conflict between ANALYZE and LISTEN/NOTIFY
E.309. Release 7.4.2

Release date: 2004-03-08

This release contains a variety of fixes from 7.4.1. For information about new features in the 7.4 major release, see Section E.311.

E.309.1. Migration to Version 7.4.2

A dump/restore is not required for those running 7.4.X. However, it might be advisable as the easiest method of incorporating fixes for two errors that have been found in the initial contents of 7.4.X system catalogs. A dump/initdb/reload sequence using 7.4.2's initdb will automatically correct these problems.

The more severe of the two errors is that data type anyarray has the wrong alignment label; this is a problem because the pg_statistic system catalog uses anyarray columns. The mislabeling can cause planner misestimations and even crashes when planning queries that involve WHERE clauses on double-aligned columns (such as float8 and timestamp). It is strongly recommended that all installations repair this error, either by initdb or by following the manual repair procedure given below.

The lesser error is that the system view pg_settings ought to be marked as having public update access, to allow UPDATE pg_settings to be used as a substitute for SET. This can also be fixed either by initdb or manually, but it is not necessary to fix unless you want to use UPDATE pg_settings.

If you wish not to do an initdb, the following procedure will work for fixing pg_statistic. As the database superuser, do:

```
-- clear out old data in pg_statistic:
DELETE FROM pg_statistic;
VACUUM pg_statistic;
-- this should update 1 row:
UPDATE pg_type SET typalign = 'd' WHERE oid = 2277;
-- this should update 6 rows:
UPDATE pg_attribute SET attalign = 'd' WHERE atttypid = 2277;
--
-- At this point you MUST start a fresh backend to avoid a crash!
--
-- repopulate pg_statistic:
ANALYZE;
```

This can be done in a live database, but beware that all backends running in the altered database must be restarted before it is safe to repopulate pg_statistic.

To repair the pg_settings error, simply do:

```
GRANT SELECT, UPDATE ON pg_settings TO PUBLIC;
```

The above procedures must be carried out in each database of an installation, including template1, and ideally including template0 as well. If you do not fix the template databases then any subsequently created databases will contain the same errors. template1 can be fixed in the same way as any other database, but fixing template0 requires additional steps. First, from any database issue:

```
UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
```
Next connect to template0 and perform the above repair procedures. Finally, do:

```
-- re-freeze template0:
VACUUM FREEZE;
-- and protect it against future alterations:
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';
```

**E.309.2. Changes**

Release 7.4.2 incorporates all the fixes included in release 7.3.6, plus the following fixes:

- Fix `pg_statistic` alignment bug that could crash optimizer
  
  See above for details about this problem.

- Allow non-super users to update `pg_settings`

- Fix several optimizer bugs, most of which led to “variable not found in subplan target lists” errors

- Avoid out-of-memory failure during startup of large multiple index scan

- Fix multibyte problem that could lead to “out of memory” error during `COPY IN`

- Fix problems with `SELECT INTO`/`CREATE TABLE AS` from tables without OIDs

- Fix problems with `alter_table` regression test during parallel testing

- Fix problems with hitting open file limit, especially on macOS (Tom)

- Partial fix for Turkish-locale issues
  
  initdb will succeed now in Turkish locale, but there are still some inconveniences associated with the i/I problem.

- Make `pg_dump` set client encoding on restore

- Other minor `pg_dump` fixes

- Allow ecpg to again use C keywords as column names (Michael)

- Added ecpg `WHENEVER NOT_FOUND` to `SELECT/INSERT/UPDATE/DELETE` (Michael)

- Fix ecpg crash for queries calling set-returning functions (Michael)

- Various other ecpg fixes (Michael)

- Fixes for Borland compiler

- Thread build improvements (Bruce)

- Various other build fixes

- Various JDBC fixes

**E.310. Release 7.4.1**

**Release date:** 2003-12-22

This release contains a variety of fixes from 7.4. For information about new features in the 7.4 major release, see Section E.311.
E.310.1. Migration to Version 7.4.1

A dump/restore is not required for those running 7.4.

If you want to install the fixes in the information schema you need to reload it into the database. This is either accomplished by initializing a new cluster by running initdb, or by running the following sequence of SQL commands in each database (ideally including template1) as a superuser in psql, after installing the new release:

```sql
DROP SCHEMA information_schema CASCADE;
\i /usr/local/pgsql/share/information_schema.sql
```

Substitute your installation path in the second command.

E.310.2. Changes

- Fixed bug in CREATE SCHEMA parsing in ECPG (Michael)
- Fix compile error when --enable-thread-safety and --with-perl are used together (Peter)
- Fix for subqueries that used hash joins (Tom)
  Certain subqueries that used hash joins would crash because of improperly shared structures.
- Fix free space map compaction bug (Tom)
  This fixes a bug where compaction of the free space map could lead to a database server shutdown.
- Fix for Borland compiler build of libpq (Bruce)
- Fix netmask() and hostmask() to return the maximum-length masklen (Tom)
  Fix these functions to return values consistent with pre-7.4 releases.
- Several contrib/pg_autovacuum fixes
  Fixes include improper variable initialization, missing vacuum after TRUNCATE, and duration computation overflow for long vacuums.
- Allow compile of contrib/cube under Cygwin (Jason Tishler)
- Fix Solaris use of password file when no passwords are defined (Tom)
  Fix crash on Solaris caused by use of any type of password authentication when no passwords were defined.
- JDBC fix for thread problems, other fixes
- Fix for bytea index lookups (Joe)
- Fix information schema for bit data types (Peter)
- Force zero_damaged_pages to be on during recovery from WAL.
- Prevent some obscure cases of “variable not in subplan target lists”
- Make PQescapeBytea and byteaout consistent with each other (Joe)
- Escape bytea output for bytes > 0x7e (Joe)
If different client encodings are used for bytea output and input, it is possible for bytea values to be corrupted by the differing encodings. This fix escapes all bytes that might be affected.

- Added missing SPI_finish() calls to dblink’s get_tuple_of_interest() (Joe)
- New Czech FAQ
- Fix information schema view constraint_column_usage for foreign keys (Peter)
- ECPG fixes (Michael)
- Fix bug with multiple IN subqueries and joins in the subqueries (Tom)
- Allow COUNT (’x’) to work (Tom)
- Install ECPG include files for Informix compatibility into separate directory (Peter)

  Some names of ECPG include files for Informix compatibility conflicted with operating system include files. By installing them in their own directory, name conflicts have been reduced.

- Fix SSL memory leak (Neil)

  This release fixes a bug in 7.4 where SSL didn’t free all memory it allocated.

- Prevent pg_service.conf from using service name as default dbname (Bruce)
- Fix local ident authentication on FreeBSD (Tom)

E.311. Release 7.4

Release date: 2003-11-17

E.311.1. Overview

Major changes in this release:

**IN*/NOT IN subqueries are now much more efficient**

In previous releases, IN/NOT IN subqueries were joined to the upper query by sequentially scanning the subquery looking for a match. The 7.4 code uses the same sophisticated techniques used by ordinary joins and so is much faster. An IN will now usually be as fast as or faster than an equivalent EXISTS subquery; this reverses the conventional wisdom that applied to previous releases.

**Improved GROUP BY processing by using hash buckets**

In previous releases, rows to be grouped had to be sorted first. The 7.4 code can do GROUP BY without sorting, by accumulating results into a hash table with one entry per group. It will still use the sort technique, however, if the hash table is estimated to be too large to fit in sort_mem.

**New multikey hash join capability**

In previous releases, hash joins could only occur on single keys. This release allows multicolumn hash joins.

**Queries using the explicit JOIN syntax are now better optimized**

Prior releases evaluated queries using the explicit JOIN syntax only in the order implied by the syntax. 7.4 allows full optimization of these queries, meaning the optimizer considers all possible
join orderings and chooses the most efficient. Outer joins, however, must still follow the declared ordering.

Faster and more powerful regular expression code

The entire regular expression module has been replaced with a new version by Henry Spencer, originally written for Tcl. The code greatly improves performance and supports several flavors of regular expressions.

Function-inlining for simple SQL functions

Simple SQL functions can now be inlined by including their SQL in the main query. This improves performance by eliminating per-call overhead. That means simple SQL functions now behave like macros.

Full support for IPv6 connections and IPv6 address data types

Previous releases allowed only IPv4 connections, and the IP data types only supported IPv4 addresses. This release adds full IPv6 support in both of these areas.

Major improvements in SSL performance and reliability

Several people very familiar with the SSL API have overhauled our SSL code to improve SSL key negotiation and error recovery.

Make free space map efficiently reuse empty index pages, and other free space management improvements

In previous releases, B-tree index pages that were left empty because of deleted rows could only be reused by rows with index values similar to the rows originally indexed on that page. In 7.4, VACUUM records empty index pages and allows them to be reused for any future index rows.

SQL-standard information schema

The information schema provides a standardized and stable way to access information about the schema objects defined in a database.

Cursors conform more closely to the SQL standard

The commands FETCH and MOVE have been overhauled to conform more closely to the SQL standard.

Cursors can exist outside transactions

These cursors are also called holdable cursors.

New client-to-server protocol

The new protocol adds error codes, more status information, faster startup, better support for binary data transmission, parameter values separated from SQL commands, prepared statements available at the protocol level, and cleaner recovery from COPY failures. The older protocol is still supported by both server and clients.

libpq and ECPG applications are now fully thread-safe

While previous libpq releases already supported threads, this release improves thread safety by fixing some non-thread-safe code that was used during database connection startup. The configure option --enable-thread-safety must be used to enable this feature.

New version of full-text indexing

A new full-text indexing suite is available in contrib/tsearch2.
New autovacuum tool

The new autovacuum tool in contrib/autovacuum monitors the database statistics tables for INSERT/UPDATE/DELETE activity and automatically vacuums tables when needed.

Array handling has been improved and moved into the server core

Many array limitations have been removed, and arrays behave more like fully-supported data types.

E.311.2. Migration to Version 7.4

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

• The server-side autocommit setting was removed and reimplemented in client applications and languages. Server-side autocommit was causing too many problems with languages and applications that wanted to control their own autocommit behavior, so autocommit was removed from the server and added to individual client APIs as appropriate.

• Error message wording has changed substantially in this release. Significant effort was invested to make the messages more consistent and user-oriented. If your applications try to detect different error conditions by parsing the error message, you are strongly encouraged to use the new error code facility instead.

• Inner joins using the explicit JOIN syntax might behave differently because they are now better optimized.

• A number of server configuration parameters have been renamed for clarity, primarily those related to logging.

• FETCH 0 or MOVE 0 now does nothing. In prior releases, FETCH 0 would fetch all remaining rows, and MOVE 0 would move to the end of the cursor.

• FETCH and MOVE now return the actual number of rows fetched/moved, or zero if at the beginning/end of the cursor. Prior releases would return the row count passed to the command, not the number of rows actually fetched or moved.

• COPY now can process files that use carriage-return or carriage-return/line-feed end-of-line sequences. Literal carriage-returns and line-feeds are no longer accepted in data values; use \r and \n instead.

• Trailing spaces are now trimmed when converting from type char(n) to varchar(n) or text. This is what most people always expected to happen anyway.

• The data type float(p) now measures p in binary digits, not decimal digits. The new behavior follows the SQL standard.

• Ambiguous date values now must match the ordering specified by the datestyle setting. In prior releases, a date specification of 10/20/03 was interpreted as a date in October even if datestyle specified that the day should be first. 7.4 will throw an error if a date specification is invalid for the current setting of datestyle.

• The functions oidrand, oidsrand, and userfntest have been removed. These functions were determined to be no longer useful.

• String literals specifying time-varying date/time values, such as 'now' or 'today' will no longer work as expected in column default expressions; they now cause the time of the table creation to be the default, not the time of the insertion. Functions such as now(), current_timestamp, or current_date should be used instead.
In previous releases, there was special code so that strings such as 'now' were interpreted at
INSERT time and not at table creation time, but this work around didn't cover all cases. Release
7.4 now requires that defaults be defined properly using functions such as now() or current-
timestamp. These will work in all situations.

- The dollar sign ($) is no longer allowed in operator names. It can instead be a non-first character
  in identifiers. This was done to improve compatibility with other database systems, and to avoid
  syntax problems when parameter placeholders ($n) are written adjacent to operators.

E.311.3. Changes

Below you will find a detailed account of the changes between release 7.4 and the previous major
release.

E.311.3.1. Server Operation Changes

- Allow IPv6 server connections (Nigel Kukard, Johan Jordaan, Bruce, Tom, Kurt Roeckx, Andrew
  Dunstan)

- Fix SSL to handle errors cleanly (Nathan Mueller)

  In prior releases, certain SSL API error reports were not handled correctly. This release fixes those
  problems.

- SSL protocol security and performance improvements (Sean Chittenden)

  SSL key renegotiation was happening too frequently, causing poor SSL performance. Also, initial
  key handling was improved.

- Print lock information when a deadlock is detected (Tom)

  This allows easier debugging of deadlock situations.

- Update /tmp socket modification times regularly to avoid their removal (Tom)

  This should help prevent /tmp directory cleaner administration scripts from removing server socket
  files.

- Enable PAM for macOS (Aaron Hillegass)

- Make B-tree indexes fully WAL-safe (Tom)

  In prior releases, under certain rare cases, a server crash could cause B-tree indexes to become
  corrupt. This release removes those last few rare cases.

- Allow B-tree index compaction and empty page reuse (Tom)

- Fix inconsistent index lookups during split of first root page (Tom)

  In prior releases, when a single-page index split into two pages, there was a brief period when
  another database session could miss seeing an index entry. This release fixes that rare failure case.

- Improve free space map allocation logic (Tom)

- Preserve free space information between server restarts (Tom)

  In prior releases, the free space map was not saved when the postmaster was stopped, so newly
  started servers had no free space information. This release saves the free space map, and reloads
  it when the server is restarted.
• Add start time to `pg_stat_activity` (Neil)
• New code to detect corrupt disk pages; erase with `zero_damaged_pages` (Tom)
• New client/server protocol: faster, no username length limit, allow clean exit from COPY (Tom)
• Add transaction status, table ID, column ID to client/server protocol (Tom)
• Add binary I/O to client/server protocol (Tom)
• Remove autocommit server setting; move to client applications (Tom)
• New error message wording, error codes, and three levels of error detail (Tom, Joe, Peter)

**E.311.3.2. Performance Improvements**

• Add hashing for `GROUP BY` aggregates (Tom)
• Make nested-loop joins be smarter about multicolumn indexes (Tom)
• Allow multikey hash joins (Tom)
• Improve constant folding (Tom)
• Add ability to inline simple SQL functions (Tom)
• Reduce memory usage for queries using complex functions (Tom)

   In prior releases, functions returning allocated memory would not free it until the query completed. This release allows the freeing of function-allocated memory when the function call completes, reducing the total memory used by functions.
• Improve GEQO optimizer performance (Tom)

   This release fixes several inefficiencies in the way the GEQO optimizer manages potential query paths.
• Allow `IN/NOT IN` to be handled via hash tables (Tom)
• Improve `NOT IN` (`subquery`) performance (Tom)
• Allow most `IN` subqueries to be processed as joins (Tom)
• Pattern matching operations can use indexes regardless of locale (Peter)

   There is no way for non-ASCII locales to use the standard indexes for LIKE comparisons. This release adds a way to create a special index for LIKE.
• Allow the postmaster to preload libraries using `preload_libraries` (Joe)

   For shared libraries that require a long time to load, this option is available so the library can be preloaded in the postmaster and inherited by all database sessions.
• Improve optimizer cost computations, particularly for subqueries (Tom)
• Avoid sort when subquery `ORDER BY` matches upper query (Tom)
• Deduce that `WHERE a.x = b.y AND b.y = 42` also means `a.x = 42` (Tom)
• Allow hash/merge joins on complex joins (Tom)
• Allow hash joins for more data types (Tom)
• Allow join optimization of explicit inner joins, disable with join_collapse_limit (Tom)
• Add parameter from_collapse_limit to control conversion of subqueries to joins (Tom)
• Use faster and more powerful regular expression code from Tcl (Henry Spencer, Tom)
• Use bit-mapped relation sets in the optimizer (Tom)
• Improve connection startup time (Tom)

The new client/server protocol requires fewer network packets to start a database session.

• Improve trigger/constraint performance (Stephan)
• Improve speed of col IN (const, const, const, ...) (Tom)
• Fix hash indexes which were broken in rare cases (Tom)
• Improve hash index concurrency and speed (Tom)

Prior releases suffered from poor hash index performance, particularly for high concurrency situations. This release fixes that, and the development group is interested in reports comparing B-tree and hash index performance.

• Align shared buffers on 32-byte boundary for copy speed improvement (Manfred Spraul)

Certain CPU's perform faster data copies when addresses are 32-byte aligned.

• Data type numeric reimplemented for better performance (Tom)

numeric used to be stored in base 100. The new code uses base 10000, for significantly better performance.

E.311.3.3. Server Configuration Changes

• Rename server parameter server_min_messages to log_min_messages (Bruce)

This was done so most parameters that control the server logs begin with log_.

• Rename show_*_stats to log_*_stats (Bruce)
• Rename show_source_port to log_source_port (Bruce)
• Rename hostname_lookup to log_hostname (Bruce)
• Add checkpoint_warning to warn of excessive checkpointing (Bruce)

In prior releases, it was difficult to determine if checkpoint was happening too frequently. This feature adds a warning to the server logs when excessive checkpointing happens.

• New read-only server parameters for localization (Tom)
• Change debug server log messages to output as DEBUG rather than LOG (Bruce)

• Prevent server log variables from being turned off by non-superusers (Bruce)

This is a security feature so non-superusers cannot disable logging that was enabled by the administrator.

• log_min_messages/client_min_messages now controls debug_* output (Bruce)

This centralizes client debug information so all debug output can be sent to either the client or server logs.
• Add macOS Rendezvous server support (Chris Campbell)

This allows macOS hosts to query the network for available PostgreSQL servers.

• Add ability to print only slow statements using log_min_duration_statement (Christopher)

This is an often requested debugging feature that allows administrators to see only slow queries in their server logs.

• Allow pg_hba.conf to accept netmasks in CIDR format (Andrew Dunstan)

This allows administrators to merge the host IP address and netmask fields into a single CIDR field in pg_hba.conf.

• New read-only parameter is_superuser (Tom)

• New parameter log_error_verbosity to control error detail (Tom)

This works with the new error reporting feature to supply additional error information like hints, file names and line numbers.

• postgres --describe-config now dumps server config variables (Aizaz Ahmed, Peter)

This option is useful for administration tools that need to know the configuration variable names and their minimums, maximums, defaults, and descriptions.

• Add new columns in pg_settings: context, type, source, min_val, max_val (Joe)

• Make default sharedBuffers 1000 and max_connections 100, if possible (Tom)

Prior versions defaulted to 64 shared buffers so PostgreSQL would start on even very old systems. This release tests the amount of shared memory allowed by the platform and selects more reasonable default values if possible. Of course, users are still encouraged to evaluate their resource load and size sharedBuffers accordingly.

• New pg_hba.conf record type hostnossl to prevent SSL connections (Jon Jensen)

In prior releases, there was no way to prevent SSL connections if both the client and server supported SSL. This option allows that capability.

• Remove parameter geqo_random_seed (Tom)

• Add server parameter regex_flavor to control regular expression processing (Tom)

• Make pg_ctl better handle nonstandard ports (Greg)

E.311.3.4. Query Changes

• New SQL-standard information schema (Peter)

• Add read-only transactions (Peter)

• Print key name and value in foreign-key violation messages (Dmitry Tkach)

• Allow users to see their own queries in pg_stat_activity (Kevin Brown)

In prior releases, only the superuser could see query strings using pg_stat_activity. Now ordinary users can see their own query strings.

• Fix aggregates in subqueries to match SQL standard (Tom)
The SQL standard says that an aggregate function appearing within a nested subquery belongs to the outer query if its argument contains only outer-query variables. Prior PostgreSQL releases did not handle this fine point correctly.

- Add option to prevent auto-addition of tables referenced in query (Nigel J. Andrews)

By default, tables mentioned in the query are automatically added to the `FROM` clause if they are not already there. This is compatible with historic POSTGRES behavior but is contrary to the SQL standard. This option allows selecting standard-compatible behavior.

- Allow `UPDATE ... SET col = DEFAULT` (Rod)

This allows `UPDATE` to set a column to its declared default value.

- Allow expressions to be used in `LIMIT/OFFSET` (Tom)

In prior releases, `LIMIT/OFFSET` could only use constants, not expressions.

- Implement `CREATE TABLE AS EXECUTE` (Neil, Peter)

**E.311.3.5. Object Manipulation Changes**

- Make `CREATE SEQUENCE` grammar more conforming to SQL:2003 (Neil)

- Add statement-level triggers (Neil)

While this allows a trigger to fire at the end of a statement, it does not allow the trigger to access all rows modified by the statement. This capability is planned for a future release.

- Add check constraints for domains (Rod)

This greatly increases the usefulness of domains by allowing them to use check constraints.

- Add `ALTER DOMAIN` (Rod)

This allows manipulation of existing domains.

- Fix several zero-column table bugs (Tom)

PostgreSQL supports zero-column tables. This fixes various bugs that occur when using such tables.

- Have `ALTER TABLE ... ADD PRIMARY KEY` add not-null constraint (Rod)

In prior releases, `ALTER TABLE ... ADD PRIMARY` would add a unique index, but not a not-null constraint. That is fixed in this release.

- Add `ALTER TABLE ... WITHOUT OIDS` (Rod)

This allows control over whether new and updated rows will have an OID column. This is most useful for saving storage space.

- Add `ALTER SEQUENCE` to modify minimum, maximum, increment, cache, cycle values (Rod)

- Add `ALTER TABLE ... CLUSTER ON` (Alvaro Herrera)

This command is used by `pg_dump` to record the cluster column for each table previously clustered. This information is used by database-wide cluster to cluster all previously clustered tables.

- Improve automatic type casting for domains (Rod, Tom)

- Allow dollar signs in identifiers, except as first character (Tom)
• Disallow dollar signs in operator names, so $x=51 works (Tom)

• Allow copying table schema using LIKE subtable, also SQL:2003 feature INCLUDING DEFAULTS (Rod)

• Add WITH GRANT OPTION clause to GRANT (Peter)
  
  This enabled GRANT to give other users the ability to grant privileges on an object.

E.311.3.6. Utility Command Changes

• Add ON COMMIT clause to CREATE TABLE for temporary tables (Gavin)
  
  This adds the ability for a table to be dropped or all rows deleted on transaction commit.

• Allow cursors outside transactions using WITH HOLD (Neil)
  
  In previous releases, cursors were removed at the end of the transaction that created them. Cursors can now be created with the WITH HOLD option, which allows them to continue to be accessed after the creating transaction has committed.

• FETCH 0 and MOVE 0 now do nothing (Bruce)
  
  In previous releases, FETCH 0 fetched all remaining rows, and MOVE 0 moved to the end of the cursor.

• Cause FETCH and MOVE to return the number of rows fetched/moved, or zero if at the beginning/end of cursor, per SQL standard (Bruce)
  
  In prior releases, the row count returned by FETCH and MOVE did not accurately reflect the number of rows processed.

• Properly handle SCROLL with cursors, or report an error (Neil)
  
  Allowing random access (both forward and backward scrolling) to some kinds of queries cannot be done without some additional work. If SCROLL is specified when the cursor is created, this additional work will be performed. Furthermore, if the cursor has been created with NO SCROLL, no random access is allowed.

• Implement SQL-compatible options FIRST, LAST, ABSOLUTE n, RELATIVE n for FETCH and MOVE (Tom)

• Allow EXPLAIN on DECLARE CURSOR (Tom)

• Allow CLUSTER to use index marked as pre-clustered by default (Alvaro Herrera)

• Allow CLUSTER to cluster all tables (Alvaro Herrera)
  
  This allows all previously clustered tables in a database to be reclustered with a single command.

• Prevent CLUSTER on partial indexes (Tom)

• Allow DOS and Mac line-endings in COPY files (Bruce)

• Disallow literal carriage return as a data value, backslash-carriage-return and \r are still allowed (Bruce)

• COPY changes (binary, \ , ) (Tom)

• Recover from COPY failure cleanly (Tom)

• Prevent possible memory leaks in COPY (Tom)
• Make TRUNCATE transaction-safe (Rod)

TRUNCATE can now be used inside a transaction. If the transaction aborts, the changes made by the TRUNCATE are automatically rolled back.

• Allow prepare/bind of utility commands like FETCH and EXPLAIN (Tom)

• Add EXPLAIN EXECUTE (Neil)

• Improve VACUUM performance on indexes by reducing WAL traffic (Tom)

• Functional indexes have been generalized into indexes on expressions (Tom)

   In prior releases, functional indexes only supported a simple function applied to one or more column names. This release allows any type of scalar expression.

• Have SHOW TRANSACTION ISOLATION match input to SET TRANSACTION ISOLATION (Tom)

• Have COMMENT ON DATABASE on nonlocal database generate a warning, rather than an error (Rod)

   Database comments are stored in database-local tables so comments on a database have to be stored in each database.

• Improve reliability of LISTEN/NOTIFY (Tom)

• Allow REINDEX to reliably reindex nonshared system catalog indexes (Tom)

   This allows system tables to be reindexed without the requirement of a standalone session, which was necessary in previous releases. The only tables that now require a standalone session for reindexing are the global system tables pg_database, pg_shadow, and pg_group.

E.311.3.7. Data Type and Function Changes

• New server parameter extra_float_digits to control precision display of floating-point numbers (Pedro Ferreira, Tom)

   This controls output precision which was causing regression testing problems.

• Allow +1300 as a numeric time-zone specifier, for FJST (Tom)

• Remove rarely used functions oidrand, oidsrand, and userfntest functions (Neil)

• Add md5() function to main server, already in contrib/pgcrypto (Joe)

   An MD5 function was frequently requested. For more complex encryption capabilities, use contrib/pgcrypto.

• Increase date range of timestamp (John Cochran)

• Change EXTRACT(EPOCH FROM timestamp) so timestamp without time zone is assumed to be in local time, not GMT (Tom)

• Trap division by zero in case the operating system doesn’t prevent it (Tom)

• Change the numeric data type internally to base 10000 (Tom)

• New hostmask() function (Greg Wickham)

• Fixes for to_char() and to_timestamp() (Karel)
• Allow functions that can take any argument data type and return any data type, using anyelement and anyarray (Joe)

This allows the creation of functions that can work with any data type.

• Arrays can now be specified as ARRAY[1,2,3].ARYAY[['a','b'], ['c','d']]. or ARRAY[ARRAY[ARRAY[2]]] (Joe)

• Allow proper comparisons for arrays, including ORDER BY and DISTINCT support (Joe)

• Allow indexes on array columns (Joe)

• Allow array concatenation with || (Joe)

• Allow WHERE qualification expr op ANY/SOME/ALL (array_expr) (Joe)

This allows arrays to behave like a list of values, for purposes like SELECT * FROM tab WHERE col IN (array_val).

• New array functions array_append, array_cat, array_lower, array_prepend, array_to_string, array_upper, string_to_array (Joe)

• Allow user defined aggregates to use polymorphic functions (Joe)

• Allow assignments to empty arrays (Joe)

• Allow 60 in seconds fields of time, timestamp, and interval input values (Tom)

Sixty-second values are needed for leap seconds.

• Allow cidr data type to be cast to text (Tom)

• Disallow invalid time zone names in SET TIME ZONE

• Trim trailing spaces when char is cast to varchar or text (Tom)

• Make float(p) measure the precision p in binary digits, not decimal digits (Tom)

• Add IPv6 support to the inet and cidr data types (Michael Graff)

• Add family() function to report whether address is IPv4 or IPv6 (Michael Graff)

• Have SHOW datestyle generate output similar to that used by SET datestyle (Tom)

• Make EXTRACT(TIMEZONE) and SET/SHOW TIME ZONE follow the SQL convention for the sign of time zone offsets, i.e., positive is east from UTC (Tom)

• Fix date_trunc('quarter', ...) (Bójthe Zoltán)

Prior releases returned an incorrect value for this function call.

• Make initcap() more compatible with Oracle (Mike Nolan)

initcap() now uppercases a letter appearing after any non-alphanumeric character, rather than only after whitespace.

• Allow only datestyle field order for date values not in ISO-8601 format (Greg)

• Add new datestyle values MDY, DMY, and YMD to set input field order; honor US and European for backward compatibility (Tom)

• String literals like 'now' or 'today' will no longer work as a column default. Use functions such as now(), current_timestamp instead. (change required for prepared statements) (Tom)
• Treat NaN as larger than any other value in \texttt{min()}/\texttt{max()} (Tom)

NaN was already sorted after ordinary numeric values for most purposes, but \texttt{min()} and \texttt{max()} didn't get this right.

• Prevent interval from suppressing :00 seconds display

• New functions \texttt{pg\_get\_triggerdef(prettyprint)} and \texttt{pg\_conversion\_is\_visible()} (Christopher)

• Allow time to be specified as 040506 or 0405 (Tom)

• Input date order must now be YYYY-MM-DD (with 4-digit year) or match \texttt{datestyle}

• Make \texttt{pg\_get\_constraintdef} support unique, primary-key, and check constraints (Christopher)

\textbf{E.311.3.8. Server-Side Language Changes}

• Prevent PL/pgSQL crash when \texttt{RETURN NEXT} is used on a zero-row record variable (Tom)

• Make PL/Python's \texttt{spi\_execute} interface handle null values properly (Andrew Bosma)

• Allow PL/pgSQL to declare variables of composite types without \texttt{\%ROWTYPE} (Tom)

• Fix PL/Python's \texttt{\_quote(\_)} function to handle big integers

• Make PL/Python an untrusted language, now called \texttt{plpythonu} (Kevin Jacobs, Tom)

The Python language no longer supports a restricted execution environment, so the trusted version of PL/Python was removed. If this situation changes, a version of PL/Python that can be used by non-superusers will be readded.

• Allow polymorphic PL/pgSQL functions (Joe, Tom)

• Allow polymorphic SQL functions (Joe)

• Improved compiled function caching mechanism in PL/pgSQL with full support for polymorphism (Joe)

• Add new parameter \texttt{\$0} in PL/pgSQL representing the function's actual return type (Joe)

• Allow PL/Tcl and PL/Python to use the same trigger on multiple tables (Tom)

• Fixed PL/Tcl's \texttt{spi\_prepare} to accept fully qualified type names in the parameter type list (Jan)

\textbf{E.311.3.9. \texttt{psql} Changes}

• Add \texttt{\pset pager always} to always use pager (Greg)

This forces the pager to be used even if the number of rows is less than the screen height. This is valuable for rows that wrap across several screen rows.

• Improve tab completion (Rod, Ross Reedstrom, Ian Barwick)

• Reorder \texttt{\?} help into groupings (Harald Armin Massa, Bruce)

• Add backslash commands for listing schemas, casts, and conversions (Christopher)

• \texttt{\encoding} now changes based on the server parameter \texttt{client\_encoding} (Tom)

In previous versions, \texttt{\encoding} was not aware of encoding changes made using \texttt{SET client\_encoding}.  

3195
• Save editor buffer into readline history (Ross)

  When \e is used to edit a query, the result is saved in the readline history for retrieval using the up arrow.

• Improve \c display (Christopher)

• Enhance HTML mode to be more standards-conforming (Greg)

• New \set AUTOCOMMIT off capability (Tom)

  This takes the place of the removed server parameter autocommit.

• New \set VERBOSITY to control error detail (Tom)

  This controls the new error reporting details.

• New prompt escape sequence %x to show transaction status (Tom)

• Long options for psql are now available on all platforms

E.311.3.10. pg_dump Changes

• Multiple pg_dump fixes, including tar format and large objects

• Allow pg_dump to dump specific schemas (Neil)

• Make pg_dump preserve column storage characteristics (Christopher)

  This preserves ALTER TABLE ... SET STORAGE information.

• Make pg_dump preserve CLUSTER characteristics (Christopher)

• Have pg_dumpall use GRANT/REVOKE to dump database-level privileges (Tom)

• Allow pg_dumpall to support the options -a, -s, -x of pg_dump (Tom)

• Prevent pg_dump from lowercasing identifiers specified on the command line (Tom)

• pg_dump options --use-set-session-authorization and --no-reconnect now do nothing, all dumps use SET SESSION AUTHORIZATION

  pg_dump no longer reconnects to switch users, but instead always uses SET SESSION AUTHORIZATION. This will reduce password prompting during restores.

• Long options for pg_dump are now available on all platforms

  PostgreSQL now includes its own long-option processing routines.

E.311.3.11. libpq Changes

• Add function PQfreemem for freeing memory on Windows, suggested for NOTIFY (Bruce)

  Windows requires that memory allocated in a library be freed by a function in the same library, hence free() doesn't work for freeing memory allocated by libpq. PQfreemem is the proper way to free libpq memory, especially on Windows, and is recommended for other platforms as well.

• Document service capability, and add sample file (Bruce)

  This allows clients to look up connection information in a central file on the client machine.

• Make PQsetdbLogin have the same defaults as PQconnectdb (Tom)
• Allow libpq to cleanly fail when result sets are too large (Tom)
• Improve performance of function PQunescapeBytea (Ben Lamb)
• Allow thread-safe libpq with configure option --enable-thread-safety (Lee Kindness, Philip Yarra)
• Allow function pqInternalNotice to accept a format string and arguments instead of just a preformatted message (Tom, Sean Chittenden)
• Control SSL negotiation with sslmode values disable, allow, prefer, and require (Jon Jensen)
• Allow new error codes and levels of text (Tom)
• Allow access to the underlying table and column of a query result (Tom)
  This is helpful for query-builder applications that want to know the underlying table and column names associated with a specific result set.
• Allow access to the current transaction status (Tom)
• Add ability to pass binary data directly to the server (Tom)
• Add function PQexecPrepared and PQsendQueryPrepared functions which perform bind/execute of previously prepared statements (Tom)

E.311.3.12. JDBC Changes
• Allow setNull on updateable result sets
• Allow executeBatch on a prepared statement (Barry)
• Support SSL connections (Barry)
• Handle schema names in result sets (Paul Sorenson)
• Add refcursor support (Nic Ferrier)

E.311.3.13. Miscellaneous Interface Changes
• Prevent possible memory leak or core dump during libpqctl shutdown (Tom)
• Add Informix compatibility to ECPG (Michael)
  This allows ECPG to process embedded C programs that were written using certain Informix extensions.
• Add type decimal to ECPG that is fixed length, for Informix (Michael)
• Allow thread-safe embedded SQL programs with configure option --enable-thread-safety (Lee Kindness, Bruce)
  This allows multiple threads to access the database at the same time.
• Moved Python client PyGreSQL to http://www.pygresql.org (Marc)

E.311.3.14. Source Code Changes
• Prevent need for separate platform geometry regression result files (Tom)
• Improved PPC locking primitive (Reinhard Max)
• New function `palloc0` to allocate and clear memory (Bruce)
• Fix locking code for s390x CPU (64-bit) (Tom)
• Allow OpenBSD to use local ident credentials (William Ahern)
• Make query plan trees read-only to executor (Tom)
• Add macOS startup scripts (David Wheeler)
• Allow libpq to compile with Borland C++ compiler (Lester Godwin, Karl Waclawek)
• Use our own version of `getopt_long()` if needed (Peter)
• Convert administration scripts to C (Peter)
• Bison >= 1.85 is now required to build the PostgreSQL grammar, if building from CVS
• Merge documentation into one book (Peter)
• Add Windows compatibility functions (Bruce)
• Allow client interfaces to compile under MinGW (Bruce)
• New `ereport()` function for error reporting (Tom)
• Support Intel compiler on Linux (Peter)
• Improve Linux startup scripts (Slawomir Sudnik, Darko Prenosil)
• Add support for AMD Opteron and Itanium (Jeffrey W. Baker, Bruce)
• Remove `--enable-recode` option from `configure`

This was no longer needed now that we have `CREATE CONVERSION`.

• Generate a compile error if spinlock code is not found (Bruce)

Platforms without spinlock code will now fail to compile, rather than silently using semaphores. This failure can be disabled with a new `configure` option.

E.311.3.15. Contrib Changes

• Change dbmirror license to BSD
• Improve earthdistance (Bruno Wolff III)
• Portability improvements to pgcrypto (Marko Kreen)
• Prevent crash in xml (John Gray, Michael Richards)
• Update oracle
• Update mysql
• Update cube (Bruno Wolff III)
• Update earthdistance to use cube (Bruno Wolff III)
• Update btree_gist (Oleg)
• New tsearch2 full-text search module (Oleg, Teodor)
• Add hash-based crosstab function to tablefuncs (Joe)
• Add serial column to order connectby() siblings in tablefuncs (Nabil Sayegh, Joe)
• Add named persistent connections to dblink (Shridhar Daithanka)
• New pg_autovacuum allows automatic VACUUM (Matthew T. O'Connor)
• Make pgbench honor environment variables PGHOST, PGPORT, PGUSER (Tatsuo)
• Improve intarray (Teodor Sigaev)
• Improve pgstattuple (Rod)
• Fix bug in metaphone() in fuzzystarmatch
• Improve adddepend (Rod)
• Update spi/timetravel (Böjthe Zoltán)
• Fix dbase -s option and improve non-ASCII handling (Thomas Behr, Márcio Smiderle)
• Remove array module because features now included by default (Joe)

E.312. Release 7.3.21

Release date: 2008-01-07

This release contains a variety of fixes from 7.3.20, including fixes for significant security issues.

This is expected to be the last PostgreSQL release in the 7.3.X series. Users are encouraged to update to a newer release branch soon.

E.312.1. Migration to Version 7.3.21

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

E.312.2. Changes

• Prevent functions in indexes from executing with the privileges of the user running VACUUM, ANALYZE, etc (Tom)

Functions used in index expressions and partial-index predicates are evaluated whenever a new table entry is made. It has long been understood that this poses a risk of trojan-horse code execution if one modifies a table owned by an untrustworthy user. (Note that triggers, defaults, check constraints, etc. pose the same type of risk.) But functions in indexes pose extra danger because they will be executed by routine maintenance operations such as VACUUM FULL, which are commonly performed automatically under a superuser account. For example, a nefarious user can execute code with superuser privileges by setting up a trojan-horse index definition and waiting for the next routine vacuum. The fix arranges for standard maintenance operations (including VACUUM, ANALYZE, REINDEX, and CLUSTER) to execute as the table owner rather than the calling user, using the same privilege-switching mechanism already used for SECURITY DEFINER functions. To prevent bypassing this security measure, execution of SET SESSION AUTHORIZATION and SET ROLE is now forbidden within a SECURITY DEFINER context. (CVE-2007-6600)

• Require non-superusers who use /contrib/dblink to use only password authentication, as a security measure (Joe)

The fix that appeared for this in 7.3.20 was incomplete, as it plugged the hole for only some dblink functions. (CVE-2007-6601, CVE-2007-3278)
• Fix potential crash in `translate()` when using a multibyte database encoding (Tom)

• Make `contrib/tablefunc's crosstab()` handle NULL rowid as a category in its own right, rather than crashing (Joe)

• Require a specific version of Autoconf to be used when re-generating the `configure` script (Peter)

This affects developers and packagers only. The change was made to prevent accidental use of untested combinations of Autoconf and PostgreSQL versions. You can remove the version check if you really want to use a different Autoconf version, but it's your responsibility whether the result works or not.

E.313. Release 7.3.20

Release date: 2007-09-17

This release contains fixes from 7.3.19.

E.313.1. Migration to Version 7.3.20

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

E.313.2. Changes

• Prevent index corruption when a transaction inserts rows and then aborts close to the end of a concurrent VACUUM on the same table (Tom)

• Make `CREATE DOMAIN ... DEFAULT NULL` work properly (Tom)

• Fix crash when `log_min_error_statement` logging runs out of memory (Tom)

• Require non-superusers who use `/contrib/dblink` to use only password authentication, as a security measure (Joe)

E.314. Release 7.3.19

Release date: 2007-04-23

This release contains fixes from 7.3.18, including a security fix.

E.314.1. Migration to Version 7.3.19

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

E.314.2. Changes

• Support explicit placement of the temporary-table schema within `search_path`, and disable searching it for functions and operators (Tom)

  This is needed to allow a security-definer function to set a truly secure value of `search_path`. Without it, an unprivileged SQL user can use temporary objects to execute code with the privileges of the security-definer function (CVE-2007-2138). See `CREATE FUNCTION` for more information.

• Fix potential-data-corruption bug in how VACUUM FULL handles UPDATE chains (Tom, Pavan Deolasee)
E.315. Release 7.3.18

Release date: 2007-02-05

This release contains a variety of fixes from 7.3.17, including a security fix.

E.315.1. Migration to Version 7.3.18

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

E.315.2. Changes

• Remove security vulnerability that allowed connected users to read backend memory (Tom)

  The vulnerability involves changing the data type of a table column used in a SQL function (CVE-2007-0555). This error can easily be exploited to cause a backend crash, and in principle might be used to read database content that the user should not be able to access.

• Fix rare bug wherein btree index page splits could fail due to choosing an infeasible split point (Heikki Linnakangas)

• Tighten security of multi-byte character processing for UTF8 sequences over three bytes long (Tom)

E.316. Release 7.3.17

Release date: 2007-01-08

This release contains a variety of fixes from 7.3.16.

E.316.1. Migration to Version 7.3.17

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

E.316.2. Changes

• to_number() and to_char(numeric) are now STABLE, not IMMUTABLE, for new initdb installs (Tom)

  This is because lc_numeric can potentially change the output of these functions.

• Improve index usage of regular expressions that use parentheses (Tom)

  This improves psql \d performance also.

E.317. Release 7.3.16

Release date: 2006-10-16

This release contains a variety of fixes from 7.3.15.

E.317.1. Migration to Version 7.3.16

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.
E.317.2. Changes

- Fix corner cases in pattern matching for psql's `\d` commands
- Fix index-corrupting bugs in `/contrib/ltrie` (Teodor)
- Back-port 7.4 spinlock code to improve performance and support 64-bit architectures better
- Fix SSL-related memory leak in libpq
- Fix backslash escaping in `/contrib/dbmirror`
- Adjust regression tests for recent changes in US DST laws

E.318. Release 7.3.15

Release date: 2006-05-23

This release contains a variety of fixes from 7.3.14, including patches for extremely serious security issues.

E.318.1. Migration to Version 7.3.15

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section E.320.

Full security against the SQL-injection attacks described in CVE-2006-2313 and CVE-2006-2314 might require changes in application code. If you have applications that embed untrustworthy strings into SQL commands, you should examine them as soon as possible to ensure that they are using recommended escaping techniques. In most cases, applications should be using subroutines provided by libraries or drivers (such as libpq's `PQescapeStringConn()`) to perform string escaping, rather than relying on ad hoc code to do it.

E.318.2. Changes

- Change the server to reject invalidly-encoded multibyte characters in all cases (Tatsuo, Tom)

  While PostgreSQL has been moving in this direction for some time, the checks are now applied uniformly to all encodings and all textual input, and are now always errors not merely warnings. This change defends against SQL-injection attacks of the type described in CVE-2006-2313.

- Reject unsafe uses of `'` in string literals

  As a server-side defense against SQL-injection attacks of the type described in CVE-2006-2314, the server now only accepts `'` and not `\` as a representation of ASCII single quote in SQL string literals. By default, `\` is rejected only when `client_encoding` is set to a client-only encoding (SJIS, BIG5, GBK, GB18030, or UHC), which is the scenario in which SQL injection is possible. A new configuration parameter `backslash_quote` is available to adjust this behavior when needed. Note that full security against CVE-2006-2314 might require client-side changes; the purpose of `backslash_quote` is in part to make it obvious that insecure clients are insecure.

- Modify libpq's string-escaping routines to be aware of encoding considerations

  This fixes libpq-using applications for the security issues described in CVE-2006-2313 and CVE-2006-2314. Applications that use multiple PostgreSQL connections concurrently should migrate to `PQescapeStringConn()` and `PQescapeByteaConn()` to ensure that escaping is done correctly for the settings in use in each database connection. Applications that do string escaping "by hand" should be modified to rely on library routines instead.
• Fix some incorrect encoding conversion functions
  \texttt{win1251\_to\_iso}, \texttt{alt\_to\_iso}, \texttt{euc\_tw\_to\_big5}, \texttt{euc\_tw\_to\_mic}, \texttt{mic\_to\_eu\_c\_tw} were all broken to varying extents.
• Clean up stray remaining uses of \texttt{\'} in strings (Bruce, Jan)
• Fix server to use custom DH SSL parameters correctly (Michael Fuhr)
• Fix various minor memory leaks

\textbf{E.319. Release 7.3.14}

\textbf{Release date:} 2006-02-14

This release contains a variety of fixes from 7.3.13.

\textbf{E.319.1. Migration to Version 7.3.14}

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.13, see Section \textbf{E.320}.

\textbf{E.319.2. Changes}

• Fix potential crash in \texttt{SET SESSION AUTHORIZATION} (CVE-2006-0553)
  An unprivileged user could crash the server process, resulting in momentary denial of service to other users, if the server has been compiled with Asserts enabled (which is not the default). Thanks to Akio Ishida for reporting this problem.
• Fix bug with row visibility logic in self-inserted rows (Tom)
  Under rare circumstances a row inserted by the current command could be seen as already valid, when it should not be. Repairs bug created in 7.3.11 release.
• Fix race condition that could lead to “file already exists” errors during pg\_clog file creation (Tom)
• Fix to allow restoring dumps that have cross-schema references to custom operators (Tom)
• Portability fix for testing presence of \texttt{finite} and \texttt{isinf} during configure (Tom)

\textbf{E.320. Release 7.3.13}

\textbf{Release date:} 2006-01-09

This release contains a variety of fixes from 7.3.12.

\textbf{E.320.1. Migration to Version 7.3.13}

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.10, see Section \textbf{E.323}. Also, you might need to \texttt{REINDEX} indexes on textual columns after updating, if you are affected by the locale or plperl issues described below.

\textbf{E.320.2. Changes}

• Fix character string comparison for locales that consider different character combinations as equal, such as Hungarian (Tom)
This might require `REINDEX` to fix existing indexes on textual columns.

- Set locale environment variables during postmaster startup to ensure that plperl won’t change the locale later

This fixes a problem that occurred if the postmaster was started with environment variables specifying a different locale than what initdb had been told. Under these conditions, any use of plperl was likely to lead to corrupt indexes. You might need `REINDEX` to fix existing indexes on textual columns if this has happened to you.

- Fix longstanding bug in `strpos()` and regular expression handling in certain rarely used Asian multi-byte character sets (Tatsuo)

- Fix bug in `/contrib/pgcrypto gen_salt`, which caused it not to use all available salt space for MD5 and XDES algorithms (Marko Kreen, Solar Designer)

  Salts for Blowfish and standard DES are unaffected.

- Fix `/contrib/dblink` to throw an error, rather than crashing, when the number of columns specified is different from what’s actually returned by the query (Joe)

### E.321. Release 7.3.12

**Release date:** 2005-12-12

This release contains a variety of fixes from 7.3.11.

#### E.321.1. Migration to Version 7.3.12

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.10, see Section E.323.

#### E.321.2. Changes

- Fix race condition in transaction log management

  There was a narrow window in which an I/O operation could be initiated for the wrong page, leading to an Assert failure or data corruption.

- `/contrib/ltree` fixes (Teodor)

- Fix longstanding planning error for outer joins

  This bug sometimes caused a bogus error “RIGHT JOIN is only supported with merge-joinable join conditions”.

- Prevent core dump in `pg_autovacuum` when a table has been dropped

### E.322. Release 7.3.11

**Release date:** 2005-10-04

This release contains a variety of fixes from 7.3.10.

#### E.322.1. Migration to Version 7.3.11

A dump/restore is not required for those running 7.3.X. However, if you are upgrading from a version earlier than 7.3.10, see Section E.323.
E.322.2. Changes

- Fix error that allowed VACUUM to remove ctid chains too soon, and add more checking in code that follows ctid links

  This fixes a long-standing problem that could cause crashes in very rare circumstances.

- Fix CHAR() to properly pad spaces to the specified length when using a multiple-byte character set (Yoshiyuki Asaba)

  In prior releases, the padding of CHAR() was incorrect because it only padded to the specified number of bytes without considering how many characters were stored.

- Fix missing rows in queries like UPDATE a=... WHERE a... with GiST index on column a

- Improve checking for partially-written WAL pages

- Improve robustness of signal handling when SSL is enabled

- Various memory leakage fixes

- Various portability improvements

- Fix PL/pgSQL to handle var := var correctly when the variable is of pass-by-reference type

E.323. Release 7.3.10

Release date: 2005-05-09

This release contains a variety of fixes from 7.3.9, including several security-related issues.

E.323.1. Migration to Version 7.3.10

A dump/restore is not required for those running 7.3.X. However, it is one possible way of handling a significant security problem that has been found in the initial contents of 7.3.X system catalogs. A dump/initdb/reload sequence using 7.3.10's initdb will automatically correct this problem.

The security problem is that the built-in character set encoding conversion functions can be invoked from SQL commands by unprivileged users, but the functions were not designed for such use and are not secure against malicious choices of arguments. The fix involves changing the declared parameter list of these functions so that they can no longer be invoked from SQL commands. (This does not affect their normal use by the encoding conversion machinery.) It is strongly recommended that all installations repair this error, either by initdb or by following the manual repair procedure given below. The error at least allows unprivileged database users to crash their server process, and might allow unprivileged users to gain the privileges of a database superuser.

If you wish not to do an initdb, perform the following procedure instead. As the database superuser, do:

BEGIN;
UPDATE pg_proc SET proargtypes[3] = 'internal'::regtype
WHERE pronamespace = 11 AND pronargs = 5
  AND proargtypes[2] = 'cstring'::regtype;
-- The command should report having updated 90 rows;
-- if not, rollback and investigate instead of committing!
COMMIT;

The above procedure must be carried out in each database of an installation, including template1, and ideally including template0 as well. If you do not fix the template databases then any subse-
quilently created databases will contain the same error. \texttt{template1} can be fixed in the same way as any other database, but fixing \texttt{template0} requires additional steps. First, from any database issue:

```
UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
```

Next connect to \texttt{template0} and perform the above repair procedure. Finally, do:

```
-- re-freeze template0:
VACUUM FREEZE;
-- and protect it against future alterations:
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';
```

### E.323.2. Changes

- Change encoding function signature to prevent misuse
- Repair ancient race condition that allowed a transaction to be seen as committed for some purposes (eg SELECT FOR UPDATE) slightly sooner than for other purposes
  
  This is an extremely serious bug since it could lead to apparent data inconsistencies being briefly visible to applications.
- Repair race condition between relation extension and VACUUM
  
  This could theoretically have caused loss of a page's worth of freshly-inserted data, although the scenario seems of very low probability. There are no known cases of it having caused more than an Assert failure.
- Fix comparisons of \texttt{TIME WITH TIME ZONE} values
  
  The comparison code was wrong in the case where the \texttt{--enable-integer-datetime\_switch} configuration switch had been used. NOTE: if you have an index on a \texttt{TIME WITH TIME ZONE} column, it will need to be \texttt{REINDEXed} after installing this update, because the fix corrects the sort order of column values.
- Fix \texttt{EXTRACT(EPOCH)} for \texttt{TIME WITH TIME ZONE} values
- Fix mis-display of negative fractional seconds in \texttt{INTERVAL} values
  
  This error only occurred when the \texttt{--enable-integer-datetime\_configuration switch} had been used.
- Additional buffer overrun checks in \texttt{plpgsql} (Neil)
- Fix \texttt{pg\_dump} to dump trigger names containing \% correctly (Neil)
- Prevent \texttt{to\_char(interval)} from dumping core for month-related formats
- Fix \texttt{contrib/pgcrypto} for newer OpenSSL builds (Marko Kreen)
- Still more 64-bit fixes for \texttt{contrib/intagg}
- Prevent incorrect optimization of functions returning \texttt{RECORD}

### E.324. Release 7.3.9

Release date: 2005-01-31
This release contains a variety of fixes from 7.3.8, including several security-related issues.

E.324.1. Migration to Version 7.3.9

A dump/restore is not required for those running 7.3.X.

E.324.2. Changes

• Disallow LOAD to non-superusers

On platforms that will automatically execute initialization functions of a shared library (this includes at least Windows and ELF-based Unixen), LOAD can be used to make the server execute arbitrary code. Thanks to NGS Software for reporting this.

• Check that creator of an aggregate function has the right to execute the specified transition functions

This oversight made it possible to bypass denial of EXECUTE permission on a function.

• Fix security and 64-bit issues in contrib/intagg

• Add needed STRICT marking to some contrib functions (Kris Jurka)

• Avoid buffer overrun when plpgsql cursor declaration has too many parameters (Neil)

• Fix planning error for FULL and RIGHT outer joins

The result of the join was mistakenly supposed to be sorted the same as the left input. This could not only deliver mis-sorted output to the user, but in case of nested merge joins could give outright wrong answers.

• Fix plperl for quote marks in tuple fields

• Fix display of negative intervals in SQL and GERMAN datestyles

E.325. Release 7.3.8

Release date: 2004-10-22

This release contains a variety of fixes from 7.3.7.

E.325.1. Migration to Version 7.3.8

A dump/restore is not required for those running 7.3.X.

E.325.2. Changes

• Repair possible failure to update hint bits on disk

Under rare circumstances this oversight could lead to “could not access transaction status” failures, which qualifies it as a potential-data-loss bug.

• Ensure that hashed outer join does not miss tuples

Very large left joins using a hash join plan could fail to output unmatched left-side rows given just the right data distribution.

• Disallow running pg_ctl as root

This is to guard against any possible security issues.
- Avoid using temp files in /tmp in make_oidjoins_check

  This has been reported as a security issue, though it's hardly worthy of concern since there is no reason for non-developers to use this script anyway.

**E.326. Release 7.3.7**

**Release date:** 2004-08-16

This release contains one critical fix over 7.3.6, and some minor items.

**E.326.1. Migration to Version 7.3.7**

A dump/restore is not required for those running 7.3.X.

**E.326.2. Changes**

- Prevent possible loss of committed transactions during crash

  Due to insufficient interlocking between transaction commit and checkpointing, it was possible for transactions committed just before the most recent checkpoint to be lost, in whole or in part, following a database crash and restart. This is a serious bug that has existed since PostgreSQL 7.1.

- Remove asymmetrical word processing in tsearch (Teodor)

- Properly schema-qualify function names when pg_dump'ing a CAST

**E.327. Release 7.3.6**

**Release date:** 2004-03-02

This release contains a variety of fixes from 7.3.5.

**E.327.1. Migration to Version 7.3.6**

A dump/restore is not required for those running 7.3.*.

**E.327.2. Changes**

- Revert erroneous changes in rule permissions checking

  A patch applied in 7.3.3 to fix a corner case in rule permissions checks turns out to have disabled rule-related permissions checks in many not-so-corner cases. This would for example allow users to insert into views they weren't supposed to have permission to insert into. We have therefore reverted the 7.3.3 patch. The original bug will be fixed in 8.0.

- Repair incorrect order of operations in GetNewTransactionId()

  This bug could result in failure under out-of-disk-space conditions, including inability to restart even after disk space is freed.

- Ensure configure selects -fno-strict-aliasing even when an external value for CFLAGS is supplied

  On some platforms, building with -fstrict-aliasing causes bugs.

- Make pg_restore handle 64-bit off_t correctly

  This bug prevented proper restoration from archive files exceeding 4 GB.
Release Notes

- Make contrib/dblink not assume that local and remote type OIDs match (Joe)
- Quote connectby()'s start_with argument properly (Joe)
- Don't crash when a rowtype argument to a plpgsql function is NULL
- Avoid generating invalid character encoding sequences in corner cases when planning LIKE operations
- Ensure text_position() cannot scan past end of source string in multibyte cases (Korea PostgreSQL Users' Group)
- Fix index optimization and selectivity estimates for LIKE operations on bytea columns (Joe)

E.328. Release 7.3.5

Release date: 2003-12-03

This has a variety of fixes from 7.3.4.

E.328.1. Migration to Version 7.3.5

A dump/restore is not required for those running 7.3.*.

E.328.2. Changes

- Force zero_damaged_pages to be on during recovery from WAL
- Prevent some obscure cases of "variable not in subplan target lists"
- Force stats processes to detach from shared memory, ensuring cleaner shutdown
- Make PQescapeBytea and byteaout consistent with each other (Joe)
- Added missing SPI_finish() calls to dblink's get_tuple_of_interest() (Joe)
- Fix for possible foreign key violation when rule rewrites INSERT (Jan)
- Support qualified type names in PL/Tcl's spi_prepare command (Jan)
- Make pg_dump handle a procedural language handler located in pg_catalog
- Make pg_dump handle cases where a custom opclass is in another schema
- Make pg_dump dump binary-compatible casts correctly (Jan)
- Fix insertion of expressions containing subqueries into rule bodies
- Fix incorrect argument processing in clusterdb script (Anand Ranganathan)
- Fix problems with dropped columns in plpython triggers
- Repair problems with to_char() reading past end of its input string (Karel)
- Fix GB18030 mapping errors (Tatsuo)
- Fix several problems with SSL error handling and asynchronous SSL I/O
- Remove ability to bind a list of values to a single parameter in JDBC (prevents possible SQL-injection attacks)
• Fix some errors in HAVE_INT64_TIMESTAMP code paths
• Fix corner case for btree search in parallel with first root page split

E.329. Release 7.3.4

Release date: 2003-07-24

This has a variety of fixes from 7.3.3.

E.329.1. Migration to Version 7.3.4

A dump/restore is not required for those running 7.3.*.

E.329.2. Changes

• Repair breakage in timestamp-to-date conversion for dates before 2000
• Prevent rare possibility of server startup failure (Tom)
• Fix bugs in interval-to-time conversion (Tom)
• Add constraint names in a few places in pg_dump (Rod)
• Improve performance of functions with many parameters (Tom)
• Fix to_ascii() buffer overruns (Tom)
• Prevent restore of database comments from throwing an error (Tom)
• Work around buggy strxfrm() present in some Solaris releases (Tom)
• Properly escape jdbc setObject() strings to improve security (Barry)

E.330. Release 7.3.3

Release date: 2003-05-22

This release contains a variety of fixes for version 7.3.2.

E.330.1. Migration to Version 7.3.3

A dump/restore is not required for those running version 7.3.*.

E.330.2. Changes

• Repair sometimes-incorrect computation of StartUpID after a crash
• Avoid slowness with lots of deferred triggers in one transaction (Stephan)
• Don’t lock referenced row when UPDATE doesn’t change foreign key’s value (Jan)
• Use -fPIC not -fpic on Sparc (Tom Callaway)
• Repair lack of schema-awareness in contrib/reindexdb
• Fix contrib/intarray error for zero-element result array (Teodor)
• Ensure createuser script will exit on control-C (Oliver)
• Fix errors when the type of a dropped column has itself been dropped
• CHECKPOINT does not cause database panic on failure in noncritical steps
• Accept 60 in seconds fields of timestamp, time, interval input values
• Issue notice, not error, if TIMESTAMP, TIME, or INTERVAL precision too large
• Fix abstime-to-time cast function (fix is not applied unless you initdb)
• Fix pg_proc entry for timestamp_izone (fix is not applied unless you initdb)
• Make EXTRACT(EPOCH FROM timestamp without time zone) treat input as local time
• 'now'::timestampz gave wrong answer if timezone changed earlier in transaction
• HAVE_INT64_TIMESTAMP code for time with timezone overwrote its input
• Accept GLOBAL TEMP/TEMPORARY as a synonym for TEMPORARY
• Avoid improper schema-privilege-check failure in foreign-key triggers
• Fix bugs in foreign-key triggers for SET DEFAULT action
• Fix incorrect time-qual check in row fetch for UPDATE and DELETE triggers
• Foreign-key clauses were parsed but ignored in ALTER TABLE ADD COLUMN
• Fix createlang script breakage for case where handler function already exists
• Fix misbehavior on zero-column tables in pg_dump, COPY, ANALYZE, other places
• Fix misbehavior of func_error() on type names containing '%'
• Fix misbehavior of replace() on strings containing '%'
• Regular-expression patterns containing certain multibyte characters failed
• Account correctly for NULLs in more cases in join size estimation
• Avoid conflict with system definition of isblank() function or macro
• Fix failure to convert large code point values in EUC_TW conversions (Tatsuo)
• Fix error recovery for SSL_read/SSL_write calls
• Don’t do early constant-folding of type coercion expressions
• Validate page header fields immediately after reading in any page
• Repair incorrect check for ungrouped variables in unnamed joins
• Fix buffer overrun in to_ascii (Guido Notari)
• contrib/ltree fixes (Teodor)
• Fix core dump in deadlock detection on machines where char is unsigned
• Avoid running out of buffers in many-way indexscan (bug introduced in 7.3)
• Fix planner’s selectivity estimation functions to handle domains properly
• Fix dbmirror memory-allocation bug (Steven Singer)
• Prevent infinite loop in \ln(numeric) due to roundoff error
• \texttt{GROUP BY} got confused if there were multiple equal \texttt{GROUP BY} items
• Fix bad plan when inherited \texttt{UPDATE/DELETE} references another inherited table
• Prevent clustering on incomplete (partial or non-NULL-storing) indexes
• Service shutdown request at proper time if it arrives while still starting up
• Fix left-links in temporary indexes (could make backwards scans miss entries)
• Fix incorrect handling of client\textunderscore encoding setting in \texttt{postgresql.conf} (Tatsuo)
• Fix failure to respond to \texttt{pg\_ctl stop -m fast} after \texttt{Async\_NotifyHandler} runs
• Fix SPI for case where rule contains multiple statements of the same type
• Fix problem with checking for wrong type of access privilege in rule query
• Fix problem with EXCEPT in \texttt{CREATE RULE}
• Prevent problem with dropping temp tables having serial columns
• Fix replace\_vars\_with\_subplan\_refs failure in complex views
• Fix regexp slowness in single-byte encodings (Tatsuo)
• Allow qualified type names in \texttt{CREATE CAST} and \texttt{DROP CAST}
• Accept \texttt{SETOF type[]} , which formerly had to be written \texttt{SETOF \_type}
• Fix \texttt{pg\_dump} core dump in some cases with procedural languages
• Force ISO datestyle in \texttt{pg\_dump} output, for portability (Oliver)
• \texttt{pg\_dump} failed to handle error return from \texttt{lo\_read} (Oleg Drokin)
• \texttt{pg\_dumpall} failed with groups having no members (Nick Eskelinen)
• \texttt{pg\_dumpall} failed to recognize \texttt{--globals-only} switch
• \texttt{pg\_restore} failed to restore blobs if \texttt{-X disable-triggers} is specified
• Repair intrafunction memory leak in \texttt{plpgsql}
• \texttt{pltcl\_s e\_log command} dumped core if given wrong parameters (Ian Harding)
• \texttt{plpython} used wrong value of \texttt{atttypmod} (Brad McLean)
• Fix improper quoting of boolean values in Python interface (D'Arcy)
• Added \texttt{addData\_Type()} method to \texttt{PGConnection} interface for JDBC
• Fixed various problems with updateable \texttt{ResultSet}s for JDBC (Shawn Green)
• Fixed various problems with \texttt{DatabaseMeta\_Data} for JDBC (Kris Jurka, Peter Royal)
• Fixed problem with parsing table ACLs in JDBC
• Better error message for character set conversion problems in JDBC
E.331. Release 7.3.2

Release date: 2003-02-04

This release contains a variety of fixes for version 7.3.1.

E.331.1. Migration to Version 7.3.2

A dump/restore is not required for those running version 7.3.*.

E.331.2. Changes

- Restore creation of OID column in CREATE TABLE AS / SELECT INTO
- Fix pg_dump core dump when dumping views having comments
- Dump DEFERRABLE/INITIALLY DEFERRED constraints properly
- Fix UPDATE when child table’s column numbering differs from parent
- Increase default value of max_fsm_relations
- Fix problem when fetching backwards in a cursor for a single-row query
- Make backward fetch work properly with cursor on SELECT DISTINCT query
- Fix problems with loading pg_dump files containing contrib/lo usage
- Fix problem with all-numeric user names
- Fix possible memory leak and core dump during disconnect in libpq
- Make plpython’s spi_execute command handle nulls properly (Andrew Bosma)
- Adjust plpython error reporting so that its regression test passes again
- Work with bison 1.875
- Handle mixed-case names properly in plpgsql’s %type (Neil)
- Fix core dump in pgtcl when executing a query rewritten by a rule
- Repair array subscript overruns (per report from Yichen Xie)
- Reduce MAX_TIME_PRECISION from 13 to 10 in floating-point case
- Correctly case-fold variable names in per-database and per-user settings
- Fix coredump in plpgsql’s RETURN NEXT when SELECT into record returns no rows
- Fix outdated use of pg_type.typprtlen in python client interface
- Correctly handle fractional seconds in timestamps in JDBC driver
- Improve performance of getImportedKeys() in JDBC
- Make shared-library symlinks work standardly on HPUX (Giles)
- Repair inconsistent rounding behavior for timestamp, time, interval
- SSL negotiation fixes (Nathan Mueller)
• Make libpq's ~/.pgpass feature work when connecting with PQconnectDB
• Update my2pg, ora2pg
• Translation updates
• Add casts between types lo and oid in contrib/lo
• fastpath code now checks for privilege to call function

E.332. Release 7.3.1

Release date: 2002-12-18

This release contains a variety of fixes for version 7.3.

E.332.1. Migration to Version 7.3.1

A dump/restore is not required for those running version 7.3. However, it should be noted that the main PostgreSQL interface library, libpq, has a new major version number for this release, which might require recompilation of client code in certain cases.

E.332.2. Changes

• Fix a core dump of COPY TO when client/server encodings don't match (Tom)
• Allow pg_dump to work with pre-7.2 servers (Philip)
• contrib/adddepend fixes (Tom)
• Fix problem with deletion of per-user/per-database config settings (Tom)
• contrib/vacuumlo fix (Tom)
• Allow 'password' encryption even when pg_shadow contains MD5 passwords (Bruce)
• contrib/dbmirror fix (Steven Singer)
• Optimizer fixes (Tom)
• contrib/tsearch fixes (Teodor Sigaev, Magnus)
• Allow locale names to be mixed case (Nicolai Tufar)
• Increment libpq library's major version number (Bruce)
• pg_hba.conf error reporting fixes (Bruce, Neil)
• Add SCO Openserver 5.0.4 as a supported platform (Bruce)
• Prevent EXPLAIN from crashing server (Tom)
• SSL fixes (Nathan Mueller)
• Prevent composite column creation via ALTER TABLE (Tom)

E.333. Release 7.3

Release date: 2002-11-27
E.333.1. Overview

Major changes in this release:

Schemas

Schemas allow users to create objects in separate namespaces, so two people or applications can have tables with the same name. There is also a public schema for shared tables. Table/index creation can be restricted by removing privileges on the public schema.

Drop Column

PostgreSQL now supports the \texttt{ALTER TABLE ... DROP COLUMN} functionality.

Table Functions

Functions returning multiple rows and/or multiple columns are now much easier to use than before. You can call such a “table function” in the \texttt{SELECT FROM} clause, treating its output like a table. Also, PL/pgSQL functions can now return sets.

Prepared Queries

PostgreSQL now supports prepared queries, for improved performance.

Dependency Tracking

PostgreSQL now records object dependencies, which allows improvements in many areas. \texttt{DROP} statements now take either \texttt{CASCADE} or \texttt{RESTRICT} to control whether dependent objects are also dropped.

Privileges

Functions and procedural languages now have privileges, and functions can be defined to run with the privileges of their creator.

Internationalization

Both multibyte and locale support are now always enabled.

Logging

A variety of logging options have been enhanced.

Interfaces

A large number of interfaces have been moved to http://gborg.postgresql.org where they can be developed and released independently.

Functions/Identifiers

By default, functions can now take up to 32 parameters, and identifiers can be up to 63 bytes long. Also, \texttt{OPAQUE} is now deprecated: there are specific “pseudo-datatypes” to represent each of the former meanings of \texttt{OPAQUE} in function argument and result types.

E.333.2. Migration to Version 7.3

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release. If your application examines the system catalogs, additional changes will be required due to the introduction of schemas in 7.3; for more information, see: http://developer.postgresql.org/~momjian/upgrade_tips_7.3.
Observe the following incompatibilities:

- Pre-6.3 clients are no longer supported.
- `pg_hba.conf` now has a column for the user name and additional features. Existing files need to be adjusted.
- Several `postgresql.conf` logging parameters have been renamed.
- `LIMIT #,#` has been disabled; use `LIMIT # OFFSET #`.
- `INSERT` statements with column lists must specify a value for each specified column. For example, `INSERT INTO tab (col1, col2) VALUES ('val1')` is now invalid. It's still allowed to supply fewer columns than expected if the `INSERT` does not have a column list.
- `serial` columns are no longer automatically `UNIQUE`; thus, an index will not automatically be created.
- A `SET` command inside an aborted transaction is now rolled back.
- `COPY` no longer considers missing trailing columns to be null. All columns need to be specified. (However, one can achieve a similar effect by specifying a column list in the `COPY` command.)
- The data type `timestamp` is now equivalent to `timestamp without time zone`, instead of `timestamp with time zone`.
- Pre-7.3 databases loaded into 7.3 will not have the new object dependencies for `serial` columns, unique constraints, and foreign keys. See the directory `contrib/adddepend/` for a detailed description and a script that will add such dependencies.
- An empty string (`''`) is no longer allowed as the input into an integer field. Formerly, it was silently interpreted as 0.

### E.333.3. Changes

#### E.333.3.1. Server Operation

- Add `pg_locks` view to show locks (Neil)
- Security fixes for password negotiation memory allocation (Neil)
- Remove support for version 0 FE/BE protocol (PostgreSQL 6.2 and earlier) (Tom)
- Reserve the last few backend slots for superusers, add parameter `superuser_reserved_connections` to control this (Nigel J. Andrews)

#### E.333.3.2. Performance

- Improve startup by calling `localtime()` only once (Tom)
- Cache system catalog information in flat files for faster startup (Tom)
- Improve caching of index information (Tom)
- Optimizer improvements (Tom, Fernando Nasser)
- Catalog caches now store failed lookups (Tom)
- Hash function improvements (Neil)
- Improve performance of query tokenization and network handling (Peter)
Release Notes

• Speed improvement for large object restore (Mario Weilguni)
• Mark expired index entries on first lookup, saving later heap fetches (Tom)
• Avoid excessive NULL bitmap padding (Manfred Koizar)
• Add BSD-licensed qsort() for Solaris, for performance (Bruce)
• Reduce per-row overhead by four bytes (Manfred Koizar)
• Fix GEOQ optimizer bug (Neil Conway)
• Make WITHOUT OID actually save four bytes per row (Manfred Koizar)
• Add default_statistics_target variable to specify ANALYZE buckets (Neil)
• Use local buffer cache for temporary tables so no WAL overhead (Tom)
• Improve free space map performance on large tables (Stephen Marshall, Tom)
• Improved WAL write concurrency (Tom)

E.333.3.3. Privileges

• Add privileges on functions and procedural languages (Peter)
• Add OWNER to CREATE DATABASE so superusers can create databases on behalf of unprivileged users (Gavin Sherry, Tom)
• Add new object privilege bits EXECUTE and USAGE (Tom)
• Add SET SESSION AUTHORIZATION DEFAULT and RESET SESSION AUTHORIZATION (Tom)
• Allow functions to be executed with the privilege of the function owner (Peter)

E.333.3.4. Server Configuration

• Server log messages now tagged with LOG, not DEBUG (Bruce)
• Add user column to pg_hba.conf (Bruce)
• Have log_connections output two lines in log file (Tom)
• Remove debug_level from postgresql.conf, now server_min_messages (Bruce)
• New ALTER DATABASE/USER ... SET command for per-user/database initialization (Peter)
• New parameters server_min_messages and client_min_messages to control which messages are sent to the server logs or client applications (Bruce)
• Allow pg_hba.conf to specify lists of users/databases separated by commas, group names prepended with +, and file names prepended with @ (Bruce)
• Remove secondary password file capability and pg_password utility (Bruce)
• Add variable db_user_namespace for database-local user names (Bruce)
• SSL improvements (Bear Giles)
• Make encryption of stored passwords the default (Bruce)
• Allow statistics collector to be reset by calling pg_stat_reset() (Christopher)
• Add log_duration parameter (Bruce)
• Rename debug_print_query to log_statement (Bruce)
• Rename show_query_stats to show_statement_stats (Bruce)
• Add param log_min_error_statement to print commands to logs on error (Gavin)

**E.333.3.5. Queries**

• Make cursors insensitive, meaning their contents do not change (Tom)
• Disable LIMIT #,# syntax; now only LIMIT # OFFSET # supported (Bruce)
• Increase identifier length to 63 (Neil, Bruce)
• UNION fixes for merging >= 3 columns of different lengths (Tom)
• Add DEFAULT key word to INSERT, e.g., INSERT ... (...) DEFAULT (...) (Rod)
• Allow views to have default values using ALTER COLUMN ... SET DEFAULT (Neil)
• Fail on INSERTs with column lists that don't supply all column values, e.g., INSERT INTO tab (col1, col2) VALUES ('val1'); (Rod)
• Fix for join aliases (Tom)
• Fix for FULL OUTER JOINs (Tom)
• Improve reporting of invalid identifier and location (Tom, Gavin)
• Fix OPEN cursor(args) (Tom)
• Allow 'ctid' to be used in a view and currtid(viewname) (Hiroshi)
• Fix for CREATE TABLE AS with UNION (Tom)
• SQL99 syntax improvements (Thomas)
• Add statement_timeout variable to cancel queries (Bruce)
• Allow prepared queries with PREPARE/EXECUTE (Neil)
• Allow FOR UPDATE to appear after LIMIT/OFFSET (Bruce)
• Add variable autocommit (Tom, David Van Wie)

**E.333.3.6. Object Manipulation**

• Make equals signs optional in CREATE DATABASE (Gavin Sherry)
• Make ALTER TABLE OWNER change index ownership too (Neil)
• New ALTER TABLE tabname ALTER COLUMN colname SET STORAGE controls TOAST storage, compression (John Gray)
• Add schema support, CREATE/DROP SCHEMA (Tom)
• Create schema for temporary tables (Tom)
• Add variable search_path for schema search (Tom)
• Add ALTER TABLE SET/DROP NOT NULL (Christopher)

• New CREATE FUNCTION volatility levels (Tom)

• Make rule names unique only per table (Tom)

• Add ‘ON tablename’ clause to DROP RULE and COMMENT ON RULE (Tom)

• Add ALTER TRIGGER RENAME (Joe)

• New current_schema() and current_schemas() inquiry functions (Tom)

• Allow functions to return multiple rows (table functions) (Joe)

• Make WITH optional in CREATE DATABASE, for consistency (Bruce)

• Add object dependency tracking (Rod, Tom)

• Add RESTRICT/CASCADE to DROP commands (Rod)

• Add ALTER TABLE DROP for non-CHECK CONSTRAINT (Rod)

• Autodestroy sequence on DROP of table with SERIAL (Rod)

• Prevent column dropping if column is used by foreign key (Rod)

• Automatically drop constraints/functions when object is dropped (Rod)

• Add CREATE/DROP OPERATOR CLASS (Bill Studenmund, Tom)

• Add ALTER TABLE DROP COLUMN (Christopher, Tom, Hiroshi)

• Prevent inherited columns from being removed or renamed (Alvaro Herrera)

• Fix foreign key constraints to not error on intermediate database states (Stephan)

• Propagate column or table renaming to foreign key constraints

• Add CREATE OR REPLACE VIEW (Gavin, Neil, Tom)

• Add CREATE OR REPLACE RULE (Gavin, Neil, Tom)

• Have rules execute alphabetically, returning more predictable values (Tom)

• Triggers are now fired in alphabetical order (Tom)

• Add /contrib/adddepend to handle pre-7.3 object dependencies (Rod)

• Allow better casting when inserting/updating values (Tom)

E.333.3.7. Utility Commands

• Have COPY TO output embedded carriage returns and newlines as \r and \n (Tom)

• Allow DELIMITER in COPY FROM to be 8-bit clean (Tatsuo)

• Make pg_dump use ALTER TABLE ADD PRIMARY KEY, for performance (Neil)

• Disable brackets in multistatement rules (Bruce)

• Disable VACUUM from being called inside a function (Bruce)

• Allow dropdb and other scripts to use identifiers with spaces (Bruce)
• Restrict database comment changes to the current database
• Allow comments on operators, independent of the underlying function (Rod)
• Rollback SET commands in aborted transactions (Tom)
• EXPLAIN now outputs as a query (Tom)
• Display condition expressions and sort keys in EXPLAIN (Tom)
• Add 'SET LOCAL var = value' to set configuration variables for a single transaction (Tom)
• Allow ANALYZE to run in a transaction (Bruce)
• Improve COPY syntax using new WITH clauses, keep backward compatibility (Bruce)
• Fix pg_dump to consistently output tags in non-ASCII dumps (Bruce)
• Make foreign key constraints clearer in dump file (Rod)
• Add COMMENT ON CONSTRAINT (Rod)
• Allow COPY TO/FROM to specify column names (Brent Verner)
• Dump UNIQUE and PRIMARY KEY constraints as ALTER TABLE (Rod)
• Have SHOW output a query result (Joe)
• Generate failure on short COPY lines rather than pad NULLs (Neil)
• Fix CLUSTER to preserve all table attributes (Alvaro Herrera)
• New pg_settings table to view/modify GUC settings (Joe)
• Add smart quoting, portability improvements to pg_dump output (Peter)
• Dump serial columns out as SERIAL (Tom)
• Enable large file support, >2G for pg_dump (Peter, Philip Warner, Bruce)
• Disallow TRUNCATE on tables that are involved in referential constraints (Rod)
• Have TRUNCATE also auto-truncate the toast table of the relation (Tom)
• Add clusterdb utility that will auto-cluster an entire database based on previous CLUSTER operations (Alvaro Herrera)
• Overhaul pg_dumpall (Peter)
• Allow REINDEX of TOAST tables (Tom)
• Implemented START TRANSACTION, per SQL99 (Neil)
• Fix rare index corruption when a page split affects bulk delete (Tom)
• Fix ALTER TABLE ... ADD COLUMN for inheritance (Alvaro Herrera)

E.333.3.8. Data Types and Functions

• Fix factorial(0) to return 1 (Bruce)
• Date/time/timezone improvements (Thomas)
• Fix for array slice extraction (Tom)
• Fix extract/date_part to report proper microseconds for timestamp (Tatsuo)
• Allow text_substr() and bytea_substr() to read TOAST values more efficiently (John Gray)
• Add domain support (Rod)
• Make WITHOUT TIME ZONE the default for TIMESTAMP and TIME data types (Thomas)
• Allow alternate storage scheme of 64-bit integers for date/time types using --enable-integer-date-times in configure (Thomas)
• Make timezone(timestamptz) return timestamp rather than a string (Thomas)
• Allow fractional seconds in date/time types for dates prior to 1BC (Thomas)
• Limit timestamp data types to 6 decimal places of precision (Thomas)
• Change timezone conversion functions from timetz() to timezone() (Thomas)
• Add configuration variables datestyle and timezone (Tom)
• Add OVERLAY(), which allows substitution of a substring in a string (Thomas)
• Add SIMILAR TO (Thomas, Tom)
• Add regular expression SUBSTRING(string FROM pat FOR escape) (Thomas)
• Add LOCALTIME and LOCALTIMESTAMP functions (Thomas)
• Add named composite types using CREATE TYPE typename AS (column) (Joe)
• Allow composite type definition in the table alias clause (Joe)
• Add new API to simplify creation of C language table functions (Joe)
• Remove ODBC-compatible empty parentheses from calls to SQL99 functions for which these parentheses do not match the standard (Thomas)
• Allow macaddr data type to accept 12 hex digits with no separators (Mike Wyer)
• Add CREATE/DROP CAST (Peter)
• Add IS DISTINCT FROM operator (Thomas)
• Add SQL99 TREAT() function, synonym for CAST() (Thomas)
• Add pg_backend_pid() to output backend pid (Bruce)
• Add IS OF / IS NOT OF type predicate (Thomas)
• Allow bit string constants without fully-specified length (Thomas)
• Allow conversion between 8-byte integers and bit strings (Thomas)
• Implement hex literal conversion to bit string literal (Thomas)
• Allow table functions to appear in the FROM clause (Joe)
• Increase maximum number of function parameters to 32 (Bruce)
• No longer automatically create index for SERIAL column (Tom)
• Add current_database() (Rod)
• Fix cash_words() to not overflow buffer (Tom)
• Add functions replace(), split_part(), to_hex() (Joe)
• Fix LIKE for bytea as a right-hand argument (Joe)
• Prevent crashes caused by SELECT cash_out(2) (Tom)
• Fix to_char(1,"FM999.99") to return a period (Karel)
• Fix trigger/type/language functions returning OPAQUE to return proper type (Tom)

E.333.3.9. Internationalization

• Add additional encodings: Korean (JOHAB), Thai (WIN874), Vietnamese (TCVN), Arabic (WIN1256), Simplified Chinese (GBK), Korean (UHC) (Eiji Tokuya)
• Enable locale support by default (Peter)
• Add locale variables (Peter)
• Escape bytes >= 0x7f for multibyte in PQescapeBytea/PQunescapeBytea (Tatsuo)
• Add locale awareness to regular expression character classes
• Enable multibyte support by default (Tatsuo)
• Add GB18030 multibyte support (Bill Huang)
• Add CREATE/DROP CONVERSION, allowing loadable encodings (Tatsuo, Kaori)
• Add pg_conversion table (Tatsuo)
• Add SQL99 CONVERT() function (Tatsuo)
• pg_dumpall, pg_controldata, and pg_resetxlog now national-language aware (Peter)
• New and updated translations

E.333.3.10. Server-side Languages

• Allow recursive SQL function (Peter)
• Change PL/Tcl build to use configured compiler and Makefile.shlib (Peter)
• Overhaul the PL/pgSQL FOUND variable to be more Oracle-compatible (Neil, Tom)
• Allow PL/pgSQL to handle quoted identifiers (Tom)
• Allow set-returning PL/pgSQL functions (Neil)
• Make PL/pgSQL schema-aware (Joe)
• Remove some memory leaks (Nigel J. Andrews, Tom)

E.333.3.11. psql

• Don’t lowercase psql \connect database name for 7.2.0 compatibility (Tom)
• Add psql \timing to time user queries (Greg Sabino Mullane)
• Have `psql \d` show index information (Greg Sabino Mullane)
• New `psql \dD` shows domains (Jonathan Eisler)
• Allow `psql` to show rules on views (Paul ?)
• Fix for `psql` variable substitution (Tom)
• Allow `psql \d` to show temporary table structure (Tom)
• Allow `psql \d` to show foreign keys (Rod)
• Fix `?` to honor `\pset pager` (Bruce)
• Have `psql` reports its version number on startup (Tom)
• Allow `\copy` to specify column names (Tom)

**E.333.3.12. libpq**

• Add `~/.pgpass` to store host/user password combinations (Alvaro Herrera)
• Add `PQunescapeBytea()` function to `libpq` (Patrick Welche)
• Fix for sending large queries over non-blocking connections (Bernhard Herzog)
• Fix for `libpq` using timers on Win9X (David Ford)
• Allow `libpq` notify to handle servers with different-length identifiers (Tom)
• Add `libpq PQescapeString()` and `PQescapeBytea()` to Windows (Bruce)
• Fix for SSL with non-blocking connections (Jack Bates)
• Add `libpq` connection timeout parameter (Denis A Ustimenko)

**E.333.3.13. JDBC**

• Allow JDBC to compile with JDK 1.4 (Dave)
• Add JDBC 3 support (Barry)
• Allows JDBC to set loglevel by adding `?loglevel=X` to the connection URL (Barry)
• Add `Driver.info()` message that prints out the version number (Barry)
• Add updateable result sets (Raghu Nidagal, Dave)
• Add support for callable statements (Paul Bethe)
• Add query cancel capability
• Add refresh row (Dave)
• Fix MD5 encryption handling for multibyte servers (Jun Kawai)
• Add support for prepared statements (Barry)

**E.333.3.14. Miscellaneous Interfaces**

• Fixed ECPG bug concerning octal numbers in single quotes (Michael)
• Move src/interfaces/libpgeasy to http://gborg.postgresql.org (Marc, Bruce)
• Improve Python interface (Elliot Lee, Andrew Johnson, Greg Copeland)
• Add libpgtcl connection close event (Gerhard Hintermayer)
• Move src/interfaces/libpq++ to http://gborg.postgresql.org (Marc, Bruce)
• Move src/interfaces/odbc to http://gborg.postgresql.org (Marc)
• Move src/interfaces/libpgeasy to http://gborg.postgresql.org (Marc, Bruce)
• Move src/interfaces/perl5 to http://gborg.postgresql.org (Marc, Bruce)
• Remove src/bin/pgaccess from main tree, now at http://www.pgaccess.org (Bruce)
• Add pg_on_connection_loss command to libpgtcl (Gerhard Hintermayer, Tom)

E.333.3.15. Source Code

• Fix for parallel make (Peter)
• AIX fixes for linking Tcl (Andreas Zeugswetter)
• Allow PL/Perl to build under Cygwin (Jason Tishler)
• Improve MIPS compiles (Peter, Oliver Elphick)
• Require Autoconf version 2.53 (Peter)
• Require readline and zlib by default in configure (Peter)
• Allow Solaris to use Intimate Shared Memory (ISM), for performance (Scott Brunza, P.J. Josh Rovero)
• Always enable syslog in compile, remove --enable-syslog option (Tatsuo)
• Always enable multibyte in compile, remove --enable-multibyte option (Tatsuo)
• Always enable locale in compile, remove --enable-locale option (Peter)
• Fix for Win9x DLL creation (Magnus Naeslund)
• Fix for link() usage by WAL code on Windows, BeOS (Jason Tishler)
• Add sys/types.h to c.h, remove from main files (Peter, Bruce)
• Fix AIX hang on SMP machines (Tomoyuki Niijima)
• AIX SMP hang fix (Tomoyuki Niijima)
• Fix pre-1970 date handling on newer glibc libraries (Tom)
• Fix PowerPC SMP locking (Tom)
• Prevent gcc -ffast-math from being used (Peter, Tom)
• Bison >= 1.50 now required for developer builds
• Kerberos 5 support now builds with Heimdal (Peter)
• Add appendix in the User's Guide which lists SQL features (Thomas)
• Improve loadable module linking to use RTLD_NOW (Tom)
• New error levels WARNING, INFO, LOG, DEBUG[1-5] (Bruce)
• New src/port directory holds replaced libc functions (Peter, Bruce)
• New pg_namespace system catalog for schemas (Tom)
• Add pg_class.relnamespace for schemas (Tom)
• Add pg_type.typnamespace for schemas (Tom)
• Add pg_proc.pronamespace for schemas (Tom)
• Restructure aggregates to have pg_proc entries (Tom)
• System relations now have their own namespace, pg_* test not required (Fernando Nasser)
• Rename TOAST index names to be *_index rather than *_idx (Neil)
• Add namespaces for operators, opclasses (Tom)
• Add additional checks to server control file (Thomas)
• New Polish FAQ (Marcin Mazurek)
• Add Posix semaphore support (Tom)
• Document need for reindex (Bruce)
• Rename some internal identifiers to simplify Windows compile (Jan, Katherine Ward)
• Add documentation on computing disk space (Bruce)
• Remove KSQO from GUC (Bruce)
• Fix memory leak in rtree (Kenneth Been)
• Modify a few error messages for consistency (Bruce)
• Remove unused system table columns (Peter)
• Make system columns NOT NULL where appropriate (Tom)
• Clean up use of sprintf in favor of snprintf() (Neil, Jukka Holappa)
• Remove OPAQUE and create specific subtypes (Tom)
• Cleanups in array internal handling (Joe, Tom)
• Disallow pg_atoi("") (Bruce)
• Remove parameter wal_files because WAL files are now recycled (Bruce)
• Add version numbers to heap pages (Tom)

E.333.3.16. Contrib

• Allow inet arrays in /contrib/array (Neil)
• GiST fixes (Teodor Sigaev, Neil)
• Upgrade /contrib/mysql
• Add /contrib/dbsize which shows table sizes without vacuum (Peter)
• Add /contrib/intagg, integer aggregator routines (mlw)
• Improve /contrib/oid2name (Neil, Bruce)
• Improve /contrib/tsearch (Oleg, Teodor Sigaev)
• Cleanups of /contrib/rserver (Alexey V. Borzov)
• Update /contrib/oracle conversion utility (Gilles Darold)
• Update /contrib/dblink (Joe)
• Improve options supported by /contrib/vacuumlo (Mario Weilguni)
• Improvements to /contrib/intarray (Oleg, Teodor Sigaev, Andrey Oktyabrski)
• Add /contrib/reindexdb utility (Shaun Thomas)
• Add indexing to /contrib/isbn_issn (Dan Weston)
• Add /contrib/dbmirror (Steven Singer)
• Improve /contrib/pgbench (Neil)
• Add /contrib/tablefunc table function examples (Joe)
• Add /contrib/ltree data type for tree structures (Teodor Sigaev, Oleg Bartunov)
• Move /contrib/pg_controldata, pg_resetxlog into main tree (Bruce)
• Fixes to /contrib/cube (Bruno Wolff)
• Improve /contrib/fulltextindex (Christopher)

E.334. Release 7.2.8

Release date: 2005-05-09

This release contains a variety of fixes from 7.2.7, including one security-related issue.

E.334.1. Migration to Version 7.2.8

A dump/restore is not required for those running 7.2.X.

E.334.2. Changes

• Repair ancient race condition that allowed a transaction to be seen as committed for some purposes (eg SELECT FOR UPDATE) slightly sooner than for other purposes

This is an extremely serious bug since it could lead to apparent data inconsistencies being briefly visible to applications.

• Repair race condition between relation extension and VACUUM

This could theoretically have caused loss of a page's worth of freshly-inserted data, although the scenario seems of very low probability. There are no known cases of it having caused more than an Assert failure.
• Fix EXTRACT(EPOCH) for TIME WITH TIME ZONE values
• Additional buffer overrun checks in plpgsql (Neil)
• Fix pg_dump to dump index names and trigger names containing % correctly (Neil)
• Prevent to_char(interval) from dumping core for month-related formats
• Fix contrib/pgcrypto for newer OpenSSL builds (Marko Kreen)

E.335. Release 7.2.7

Release date: 2005-01-31

This release contains a variety of fixes from 7.2.6, including several security-related issues.

E.335.1. Migration to Version 7.2.7

A dump/restore is not required for those running 7.2.X.

E.335.2. Changes

• Disallow LOAD to non-superusers

  On platforms that will automatically execute initialization functions of a shared library (this includes at least Windows and ELF-based Unixen), LOAD can be used to make the server execute arbitrary code. Thanks to NGS Software for reporting this.

• Add needed STRICT marking to some contrib functions (Kris Jurka)

• Avoid buffer overrun when plpgsql cursor declaration has too many parameters (Neil)

• Fix planning error for FULL and RIGHT outer joins

  The result of the join was mistakenly supposed to be sorted the same as the left input. This could not only deliver mis-sorted output to the user, but in case of nested merge joins could give outright wrong answers.

• Fix display of negative intervals in SQL and GERMAN datestyles

E.336. Release 7.2.6

Release date: 2004-10-22

This release contains a variety of fixes from 7.2.5.

E.336.1. Migration to Version 7.2.6

A dump/restore is not required for those running 7.2.X.

E.336.2. Changes

• Repair possible failure to update hint bits on disk

  Under rare circumstances this oversight could lead to “could not access transaction status” failures, which qualifies it as a potential-data-loss bug.

• Ensure that hashed outer join does not miss tuples
Very large left joins using a hash join plan could fail to output unmatched left-side rows given just the right data distribution.

- Disallow running pg_ctl as root
  This is to guard against any possible security issues.
- Avoid using temp files in /tmp in make_oidjoins_check
  This has been reported as a security issue, though it's hardly worthy of concern since there is no reason for non-developers to use this script anyway.
- Update to newer versions of Bison

E.337. Release 7.2.5

Release date: 2004-08-16

This release contains a variety of fixes from 7.2.4.

E.337.1. Migration to Version 7.2.5

A dump/restore is not required for those running 7.2.X.

E.337.2. Changes

- Prevent possible loss of committed transactions during crash
  Due to insufficient interlocking between transaction commit and checkpointing, it was possible for transactions committed just before the most recent checkpoint to be lost, in whole or in part, following a database crash and restart. This is a serious bug that has existed since PostgreSQL 7.1.
- Fix corner case for btree search in parallel with first root page split
- Fix buffer overrun in to_ascii (Guido Notari)
- Fix core dump in deadlock detection on machines where char is unsigned
- Fix failure to respond to pg_ctl stop -m fast after Async_NotifyHandler runs
- Repair memory leaks in pg_dump
- Avoid conflict with system definition of isblank() function or macro

E.338. Release 7.2.4

Release date: 2003-01-30

This release contains a variety of fixes for version 7.2.3, including fixes to prevent possible data loss.

E.338.1. Migration to Version 7.2.4

A dump/restore is not required for those running version 7.2.*.

E.338.2. Changes

- Fix some additional cases of VACUUM "No one parent tuple was found" error
• Prevent VACUUM from being called inside a function (Bruce)
• Ensure pg_clog updates are sync'd to disk before marking checkpoint complete
• Avoid integer overflow during large hash joins
• Make GROUP commands work when pg_group.grolist is large enough to be toasted
• Fix errors in datetime tables; some timezone names weren't being recognized
• Fix integer overflows in circle_poly(), path_encode(), path_add() (Neil)
• Repair long-standing logic errors in lseg_eq(), lseg_ne(), lseg_center()

E.339. Release 7.2.3

Release date: 2002-10-01

This release contains a variety of fixes for version 7.2.2, including fixes to prevent possible data loss.

E.339.1. Migration to Version 7.2.3

A dump/restore is not required for those running version 7.2.*.

E.339.2. Changes

• Prevent possible compressed transaction log loss (Tom)
• Prevent non-superuser from increasing most recent vacuum info (Tom)
• Handle pre-1970 date values in newer versions of glibc (Tom)
• Fix possible hang during server shutdown
• Prevent spinlock hangs on SMP PPC machines (Tomoyuki Niijima)
• Fix pg_dump to properly dump FULL JOIN USING (Tom)

E.340. Release 7.2.2

Release date: 2002-08-23

This release contains a variety of fixes for version 7.2.1.

E.340.1. Migration to Version 7.2.2

A dump/restore is not required for those running version 7.2.*.

E.340.2. Changes

• Allow EXECUTE of "CREATE TABLE AS ... SELECT" in PL/pgSQL (Tom)
• Fix for compressed transaction log id wraparound (Tom)
• Fix PQescapeBytea/PQunescapeBytea so that they handle bytes > 0x7f (Tatsuo)
• Fix psql and pg_dump crashing when invoked with non-existent long options (Tatsuo)
Release Notes

• Fix crash when invoking geometric operators (Tom)
• Allow OPEN cursor(args) (Tom)
• Fix for rtree_gist index build (Teodor)
• Fix for dumping user-defined aggregates (Tom)
• contrib/intarray fixes (Oleg)
• Fix for complex UNION/EXCEPT/INTERSECT queries using parens (Tom)
• Fix to pg_convert (Tatsuo)
• Fix for crash with long DATA strings (Thomas, Neil)
• Fix for repeat(), lpad(), rpad() and long strings (Neil)

E.341. Release 7.2.1

Release date: 2002-03-21

This release contains a variety of fixes for version 7.2.

E.341.1. Migration to Version 7.2.1

A dump/restore is not required for those running version 7.2.

E.341.2. Changes

• Ensure that sequence counters do not go backwards after a crash (Tom)
• Fix pgaccess kanji-conversion key binding (Tatsuo)
• Optimizer improvements (Tom)
• Cash I/O improvements (Tom)
• New Russian FAQ
• Compile fix for missing AuthBlockSig (Heiko)
• Additional time zones and time zone fixes (Thomas)
• Allow psql \connect to handle mixed case database and user names (Tom)
• Return proper OID on command completion even with ON INSERT rules (Tom)
• Allow COPY FROM to use 8-bit DELIMITERS (Tatsuo)
• Fix bug in extract/date_part for milliseconds/microseconds (Tatsuo)
• Improve handling of multiple UNIONs with different lengths (Tom)
• contrib/btree_gist improvements (Teodor Sigaev)
• contrib/tsearch dictionary improvements, see README.tsearch for an additional installation step (Thomas T. Thai, Teodor Sigaev)
• Fix for array subscripts handling (Tom)
• Allow EXECUTE of "CREATE TABLE AS ... SELECT" in PL/pgSQL (Tom)

E.342. Release 7.2

Release date: 2002-02-04

E.342.1. Overview

This release improves PostgreSQL for use in high-volume applications.

Major changes in this release:

VACUUM

Vacuuming no longer locks tables, thus allowing normal user access during the vacuum. A new VACUUM FULL command does old-style vacuum by locking the table and shrinking the on-disk copy of the table.

Transactions

There is no longer a problem with installations that exceed four billion transactions.

OIDs

OIDs are now optional. Users can now create tables without OIDs for cases where OID usage is excessive.

Optimizer

The system now computes histogram column statistics during ANALYZE, allowing much better optimizer choices.

Security

A new MD5 encryption option allows more secure storage and transfer of passwords. A new Unix-domain socket authentication option is available on Linux and BSD systems.

Statistics

Administrators can use the new table access statistics module to get fine-grained information about table and index usage.

Internationalization

Program and library messages can now be displayed in several languages.

E.342.2. Migration to Version 7.2

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

Observe the following incompatibilities:

• The semantics of the VACUUM command have changed in this release. You might wish to update your maintenance procedures accordingly.

• In this release, comparisons using = NULL will always return false (or NULL, more precisely). Previous releases automatically transformed this syntax to IS NULL. The old behavior can be re-enabled using a postgresql.conf parameter.
• The `pg_hba.conf` and `pg_ident.conf` configuration is now only reloaded after receiving a SIGHUP signal, not with each connection.

• The function `octet_length()` now returns the uncompressed data length.

• The date/time value 'current' is no longer available. You will need to rewrite your applications.

• The `timestamp()`, `time()`, and `interval()` functions are no longer available. Instead of `timestamp()`, use `timestamp 'string'` or `CAST`.

The `SELECT ... LIMIT #,#` syntax will be removed in the next release. You should change your queries to use separate LIMIT and OFFSET clauses, e.g. `LIMIT 10 OFFSET 20`.

**E.342.3. Changes**

**E.342.3.1. Server Operation**

• Create temporary files in a separate directory (Bruce)

• Delete orphaned temporary files on postmaster startup (Bruce)

• Added unique indexes to some system tables (Tom)

• System table operator reorganization (Oleg Bartunov, Teodor Sigaev, Tom)

• Renamed `pg_log` to `pg_clog` (Tom)

• Enable SIGTERM, SIGQUIT to kill backends (Jan)

• Removed compile-time limit on number of backends (Tom)

• Better cleanup for semaphore resource failure (Tatsuo, Tom)

• Allow safe transaction ID wraparound (Tom)

• Removed OIDs from some system tables (Tom)

• Removed "triggered data change violation" error check (Tom)

• SPI portal creation of prepared/saved plans (Jan)

• Allow SPI column functions to work for system columns (Tom)

• Long value compression improvement (Tom)

• Statistics collector for table, index access (Jan)

• Truncate extra-long sequence names to a reasonable value (Tom)

• Measure transaction times in milliseconds (Thomas)

• Fix TID sequential scans (Hiroshi)

• Superuser ID now fixed at 1 (Peter E)

• New `pg_ctl "reload"` option (Tom)

**E.342.3.2. Performance**

• Optimizer improvements (Tom)
• New histogram column statistics for optimizer (Tom)
• Reuse write-ahead log files rather than discarding them (Tom)
• Cache improvements (Tom)
• IS NULL, IS NOT NULL optimizer improvement (Tom)
• Improve lock manager to reduce lock contention (Tom)
• Keep relcache entries for index access support functions (Tom)
• Allow better selectivity with NaN and infinities in NUMERIC (Tom)
• R-tree performance improvements (Kenneth Been)
• B-tree splits more efficient (Tom)

E.342.3.3. Privileges

• Change UPDATE, DELETE privileges to be distinct (Peter E)
• New REFERENCES, TRIGGER privileges (Peter E)
• Allow GRANT/REVOKE to/from more than one user at a time (Peter E)
• New has_table_privilege() function (Joe Conway)
• Allow non-superuser to vacuum database (Tom)
• New SET SESSION AUTHORIZATION command (Peter E)
• Fix bug in privilege modifications on newly created tables (Tom)
• Disallow access to pg_statistic for non-superuser, add user-accessible views (Tom)

E.342.3.4. Client Authentication

• Fork postmaster before doing authentication to prevent hangs (Peter E)
• Add ident authentication over Unix domain sockets on Linux, *BSD (Helge Bahmann, Oliver Elphick, Teodor Sigaev, Bruce)
• Add a password authentication method that uses MD5 encryption (Bruce)
• Allow encryption of stored passwords using MD5 (Bruce)
• PAM authentication (Dominic J. Eidson)
• Load pg_hba.conf and pg_ident.conf only on startup and SIGHUP (Bruce)

E.342.3.5. Server Configuration

• Interpretation of some time zone abbreviations as Australian rather than North American now settable at run time (Bruce)
• New parameter to set default transaction isolation level (Peter E)
• New parameter to enable conversion of "expr = NULL" into "expr IS NULL", off by default (Peter E)
• New parameter to control memory usage by VACUUM (Tom)
• New parameter to set client authentication timeout (Tom)
• New parameter to set maximum number of open files (Tom)

E.342.3.6. Queries

• Statements added by INSERT rules now execute after the INSERT (Jan)
• Prevent unadorned relation names in target list (Bruce)
• NULLs now sort after all normal values in ORDER BY (Tom)
• New IS UNKNOWN, IS NOT UNKNOWN Boolean tests (Tom)
• New SHARE UPDATE EXCLUSIVE lock mode (Tom)
• New EXPLAIN ANALYZE command that shows run times and row counts (Martijn van Oosterhout)
• Fix problem with LIMIT and subqueries (Tom)
• Fix for LIMIT, DISTINCT ON pushed into subqueries (Tom)
• Fix nested EXCEPT/INTERSECT (Tom)

E.342.3.7. Schema Manipulation

• Fix SERIAL in temporary tables (Bruce)
• Allow temporary sequences (Bruce)
• Sequences now use int8 internally (Tom)
• New SERIAL8 creates int8 columns with sequences, default still SERIAL4 (Tom)
• Make OIDs optional using WITHOUT OIDS (Tom)
• Add %TYPE syntax to CREATE TYPE (Ian Lance Taylor)
• Add ALTER TABLE / DROP CONSTRAINT for CHECK constraints (Christopher Kings-Lynne)
• New CREATE OR REPLACE FUNCTION to alter existing function (preserving the function OID) (Gavin Sherry)
• Add ALTER TABLE / ADD [ UNIQUE | PRIMARY ] (Christopher Kings-Lynne)
• Allow column renaming in views
• Make ALTER TABLE / RENAME COLUMN update column names of indexes (Brent Verner)
• Fix for ALTER TABLE / ADD CONSTRAINT ... CHECK with inherited tables (Stephan Szabo)
• ALTER TABLE RENAME update foreign-key trigger arguments correctly (Brent Verner)
• DROP AGGREGATE and COMMENT ON AGGREGATE now accept an aggtype (Tom)
• Add automatic return type data casting for SQL functions (Tom)
• Allow GiST indexes to handle NULLs and multikey indexes (Oleg Bartunov, Teodor Sigaev, Tom)
• Enable partial indexes (Martijn van Oosterhout)

E.342.3.8. Utility Commands

• Add RESET ALL, SHOW ALL (Marko Kreen)
• CREATE/ALTER USER/GROUP now allow options in any order (Vince)
• Add LOCK A, B, C functionality (Neil Padgett)
• New ENCRYPTED/UNENCRYPTED option to CREATE/ALTER USER (Bruce)
• New light-weight VACUUM does not lock table; old semantics are available as VACUUM FULL (Tom)
• Disable COPY TO/FROM on views (Bruce)
• COPY DELIMITERS string must be exactly one character (Tom)
• VACUUM warning about index tuples fewer than heap now only appears when appropriate (Martijn van Oosterhout)
• Fix privilege checks for CREATE INDEX (Tom)
• Disallow inappropriate use of CREATE/DROP INDEX/TRIGGER/VIEWS (Tom)

E.342.3.9. Data Types and Functions

• SUM(), AVG(), COUNT() now uses int8 internally for speed (Tom)
• Add convert(), convert2() (Tatsuo)
• New function bit_length() (Peter E)
• Make the "n" in CHAR(n)/VARCHAR(n) represents letters, not bytes (Tatsuo)
• CHAR(), VARCHAR() now reject strings that are too long (Peter E)
• BIT VARYING now rejects bit strings that are too long (Peter E)
• BIT now rejects bit strings that do not match declared size (Peter E)
• INET, CIDR text conversion functions (Alex Pilosov)
• INET, CIDR operators << and <<= indexable (Alex Pilosov)
• Bytea \### now requires valid three digit octal number
• Bytea comparison improvements, now supports =, <>, >, >=, <, and <=
• Bytea now supports B-tree indexes
• Bytea now supports LIKE, LIKE...ESCAPE, NOT LIKE, NOT LIKE...ESCAPE
• Bytea now supports concatenation
• New bytea functions: position, substring, trim, btrim, and length
• New encode() function mode, "escaped", converts minimally escaped bytea to/from text
• Add pg_database_encoding_max_length() (Tatsuo)
• Add pg_client_encoding() function (Tatsuo)
• now() returns time with millisecond precision (Thomas)
• New TIMESTAMP WITHOUT TIMEZONE data type (Thomas)
• Add ISO date/time specification with "T", yyyy-mm-ddThh:mm:ss (Thomas)
• New xid/int comparison functions (Hiroshi)
• Add precision to TIME, TIMESTAMP, and INTERVAL data types (Thomas)
• Modify type coercion logic to attempt binary-compatible functions first (Tom)
• New encode() function installed by default (Marko Kreen)
• Improved to_() conversion functions (Karel Zak)
• Optimize LIKE/ILIKE when using single-byte encodings (Tatsuo)
• New functions in contrib/pgcrypto: crypt(), hmac(), encrypt(), gen_salt() (Marko Kreen)
• Correct description of translate() function (Bruce)
• Add INTERVAL argument for SET TIME ZONE (Thomas)
• Add INTERVAL YEAR TO MONTH (etc.) syntax (Thomas)
• Optimize length functions when using single-byte encodings (Tatsuo)
• Fix path_inter, path_distance, path_length, dist_ppath to handle closed paths (Curtis Barrett, Tom)
• octet_length(text) now returns non-compressed length (Tatsuo, Bruce)
• Handle "July" full name in date/time literals (Greg Sabino Mullane)
• Some datatype() function calls now evaluated differently
• Add support for Julian and ISO time specifications (Thomas)

E.342.3.10. Internationalization

• National language support in psql, pg_dump, libpq, and server (Peter E)
• Message translations in Chinese (simplified, traditional), Czech, French, German, Hungarian, Russian, Swedish (Peter E, Serguei A. Mokhov, Karel Zak, Weiping He, Zhenbang Wei, Kovacs Zoltan)
• Make trim, ltrim, rtrim, btrim, lpad, rpad, translate multibyte aware (Tatsuo)
• Add LATIN5,6,7,8,9,10 support (Tatsuo)
• Add ISO 8859-5,6,7,8,9,10 support (Tatsuo)
• Correct LATIN5 to mean ISO-8859-9, not ISO-8859-5 (Tatsuo)
• Make mic2ascii() non-ASCII aware (Tatsuo)
• Reject invalid multibyte character sequences (Tatsuo)

E.342.3.11. PL/pgSQL

• Now uses portals for SELECT loops, allowing huge result sets (Jan)

3236
• CURSOR and REFCURSOR support (Jan)
• Can now return open cursors (Jan)
• Add ELSEIF (Klaus Reger)
• Improve PL/pgSQL error reporting, including location of error (Tom)
• Allow IS or FOR key words in cursor declaration, for compatibility (Bruce)
• Fix for SELECT ... FOR UPDATE (Tom)
• Fix for PERFORM returning multiple rows (Tom)
• Make PL/pgSQL use the server's type coercion code (Tom)
• Memory leak fix (Jan, Tom)
• Make trailing semicolon optional (Tom)

E.342.3.12. PL/Perl
• New untrusted PL/Perl (Alex Pilosov)
  • PL/Perl is now built on some platforms even if libperl is not shared (Peter E)

E.342.3.13. PL/Tcl
• Now reports errorInfo (Vsevolod Lobko)
  • Add spi_lastoid function (bob@redivi.com)

E.342.3.14. PL/Python
  • ...is new (Andrew Bosma)

E.342.3.15. psql
• \d displays indexes in unique, primary groupings (Christopher Kings-Lynne)
• Allow trailing semicolons in backslash commands (Greg Sabino Mullane)
• Read password from /dev/tty if possible
• Force new password prompt when changing user and database (Tatsuo, Tom)
• Format the correct number of columns for Unicode (Patrice)

E.342.3.16. libpq
• New function PQescapeString() to escape quotes in command strings (Florian Weimer)
  • New function PQescapeBytea() escapes binary strings for use as SQL string literals

E.342.3.17. JDBC
• Return OID of INSERT (Ken K)
• Handle more data types (Ken K)
• Handle single quotes and newlines in strings (Ken K)
• Handle NULL variables (Ken K)
• Fix for time zone handling (Barry Lind)
• Improved Druid support
• Allow eight-bit characters with non-multibyte server (Barry Lind)
• Support BIT, BINARY types (Ned Wolpert)
• Reduce memory usage (Michael Stephens, Dave Cramer)
• Update DatabaseMetaData (Peter E)
• Add DatabaseMetaData.getCatalogs() (Peter E)
• Encoding fixes (Anders Bengtsson)
• Get/setCatalog methods (Jason Davies)
• DatabaseMetaData.getColumn() now returns column defaults (Jason Davies)
• DatabaseMetaData.getColumn() performance improvement (Jeroen van Vianen)
• Some JDBC1 and JDBC2 merging (Anders Bengtsson)
• Transaction performance improvements (Barry Lind)
• Array fixes (Greg Zoller)
• Serialize addition
• Fix batch processing (Rene Pijlman)
• ExecSQL method reorganization (Anders Bengtsson)
• GetColumn() fixes (Jeroen van Vianen)
• Fix isWriteable() function (Rene Pijlman)
• Improved passage of JDBC2 conformance tests (Rene Pijlman)
• Add bytea type capability (Barry Lind)
• Add isNullable() (Rene Pijlman)
• JDBC date/time test suite fixes (Liam Stewart)
• Fix for SELECT `id` AS xxx FROM table (Dave Cramer)
• Fix DatabaseMetaData to show precision properly (Mark Lillywhite)
• New getImported/getExported keys (Jason Davies)
• MD5 password encryption support (Jeremy Wohl)
• Fix to actually use type cache (Ned Wolpert)

E.342.3.18. ODBC

• Remove query size limit (Hiroshi)
Release Notes

- Remove text field size limit (Hiroshi)
- Fix for SQLPrimaryKeys in multibyte mode (Hiroshi)
- Allow ODBC procedure calls (Hiroshi)
- Improve boolean handing (Aidan Mountford)
- Most configuration options now settable via DSN (Hiroshi)
- Multibyte, performance fixes (Hiroshi)
- Allow driver to be used with iODBC or unixODBC (Peter E)
- MD5 password encryption support (Bruce)
- Add more compatibility functions to odbc.sql (Peter E)

E.342.3.19. ECPG

- EXECUTE ... INTO implemented (Christof Petig)
- Multiple row descriptor support (e.g. CARDINALITY) (Christof Petig)
- Fix for GRANT parameters (Lee Kindness)
- Fix INITIALLY DEFERRED bug
- Various bug fixes (Michael, Christof Petig)
- Auto allocation for indicator variable arrays (int *ind_p=NULL)
- Auto allocation for string arrays (char **foo_pp=NULL)
- ECPGfree_auto_mem fixed
- All function names with external linkage are now prefixed by ECPG
- Fixes for arrays of structures (Michael)

E.342.3.20. Misc. Interfaces

- Python fix fetchone() (Gerhard Haring)
- Use UTF, Unicode in Tcl where appropriate (Vsevolod Lobko, Reinhard Max)
- Add Tcl COPY TO/FROM (ljb)
- Prevent output of default index op class in pg_dump (Tom)
- Fix libpgeasy memory leak (Bruce)

E.342.3.21. Build and Install

- Configure, dynamic loader, and shared library fixes (Peter E)
- Fixes in QNX 4 port (Bernd Tegge)
- Fixes in Cygwin and Windows ports (Jason Tishler, Gerhard Haring, Dmitry Yurtaev, Darko Prenosil, Mikhail Terekhov)
• Fix for Windows socket communication failures (Magnus, Mikhail Terekhov)
• Hurd compile fix (Oliver Elphick)
• BeOS fixes (Cyril Velter)
• Remove configure --enable-unicode-conversion, now enabled by multibyte (Tatsuo)
• AIX fixes (Tatsuo, Andreas)
• Fix parallel make (Peter E)
• Install SQL language manual pages into OS-specific directories (Peter E)
• Rename config.h to pg_config.h (Peter E)
• Reorganize installation layout of header files (Peter E)

E.342.3.22. Source Code

• Remove SEP_CHAR (Bruce)
• New GUC hooks (Tom)
• Merge GUC and command line handling (Marko Kreen)
• Remove EXTEND INDEX (Martijn van Oosterhout, Tom)
• New pgjindent utility to indent java code (Bruce)
• Remove define of true/false when compiling under C++ (Leandro Fanzone, Tom)
• pgindent fixes (Bruce, Tom)
• Replace strcasecmp() with strcmp() where appropriate (Peter E)
• Dynahash portability improvements (Tom)
• Add 'volatile' usage in spinlock structures
• Improve signal handling logic (Tom)

E.342.3.23. Contrib

• New contrib/rtree_gist (Oleg Bartunov, Teodor Sigaev)
• New contrib/tsearch full-text indexing (Oleg, Teodor Sigaev)
• Add contrib/dblink for remote database access (Joe Conway)
• contrib/ora2pg Oracle conversion utility (Gilles Darold)
• contrib/xml XML conversion utility (John Gray)
• contrib/fulltextindex fixes (Christopher Kings-Lynne)
• New contrib/fuzzystrmatch with levenshtein and metaphone, soundex merged (Joe Conway)
• Add contrib/intarray boolean queries, binary search, fixes (Oleg Bartunov)
• New pg_upgrade utility (Bruce)
• Add new pg_resetxlog options (Bruce, Tom)

E.343. Release 7.1.3

Release date:  2001-08-15

E.343.1. Migration to Version 7.1.3

A dump/restore is not required for those running 7.1.X.

E.343.2. Changes

Remove unused WAL segments of large transactions (Tom)
Multiactio rule fix (Tom)
PL/pgSQL memory allocation fix (Jan)
VACUUM buffer fix (Tom)
Regression test fixes (Tom)
pg_dump fixes for GRANT/REVOKE/comments on views, user-defined types (Tom)
Fix subselects with DISTINCT ON or LIMIT (Tom)
BeOS fix
Disable COPY TO/FROM a view (Tom)
Cygwin build (Jason Tishler)

E.344. Release 7.1.2

Release date:  2001-05-11

This has one fix from 7.1.1.

E.344.1. Migration to Version 7.1.2

A dump/restore is not required for those running 7.1.X.

E.344.2. Changes

Fix PL/pgSQL SELECTs when returning no rows
Fix for psql backslash core dump
Referential integrity privilege fix
Optimizer fixes
pg_dump cleanups

E.345. Release 7.1.1

Release date:  2001-05-05

This has a variety of fixes from 7.1.

E.345.1. Migration to Version 7.1.1

A dump/restore is not required for those running 7.1.
E.345.2. Changes

Fix for numeric MODULO operator (Tom)
pg_dump fixes (Philip)
pg_dump can dump 7.0 databases (Philip)
readline 4.2 fixes (Peter E)
JOIN fixes (Tom)
AIX, MSWIN, VAX, N32K fixes (Tom)
Multibytes fixes (Tom)
Unicode fixes (Tatsuo)
Optimizer improvements (Tom)
Fix for whole rows in functions (Tom)
Fix for pg_ctl and option strings with spaces (Peter E)
ODBC fixes (Hiroshi)
EXTRACT can now take string argument (Thomas)
Python fixes (Darcy)

E.346. Release 7.1

Release date: 2001-04-13

This release focuses on removing limitations that have existed in the PostgreSQL code for many years.

Major changes in this release:

Write-ahead Log (WAL)

To maintain database consistency in case of an operating system crash, previous releases of PostgreSQL have forced all data modifications to disk before each transaction commit. With WAL, only one log file must be flushed to disk, greatly improving performance. If you have been using -F in previous releases to disable disk flushes, you might want to consider discontinuing its use.

TOAST

TOAST - Previous releases had a compiled-in row length limit, typically 8k - 32k. This limit made storage of long text fields difficult. With TOAST, long rows of any length can be stored with good performance.

Outer Joins

We now support outer joins. The UNION/NOT IN workaround for outer joins is no longer required. We use the SQL92 outer join syntax.

Function Manager

The previous C function manager did not handle null values properly, nor did it support 64-bit CPU's (Alpha). The new function manager does. You can continue using your old custom functions, but you might want to rewrite them in the future to use the new function manager call interface.

Complex Queries

A large number of complex queries that were unsupported in previous releases now work. Many combinations of views, aggregates, UNION, LIMIT, cursors, subqueries, and inherited tables now work properly. Inherited tables are now accessed by default. Subqueries in FROM are now supported.
E.346.1. Migration to Version 7.1

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release.

E.346.2. Changes

Bug Fixes

--------

Many multibyte/Unicode/locale fixes (Tatsuo and others)
More reliable ALTER TABLE RENAME (Tom)
Kerberos V fixes (David Wragg)
Fix for INSERT INTO...SELECT where targetlist has subqueries (Tom)
Prompt username/password on standard error (Bruce)
Large objects inv_read/inv_write fixes (Tom)
Fixes for to_char(), to_date(), to_ascii(), and to_timestamp()
(Karel, Daniel Baldoni)
Prevent query expressions from leaking memory (Tom)
Allow UPDATE of arrays elements (Tom)
Wake up lock waiters during cancel (Hiroshi)
Fix rare cursor crash when using hash join (Tom)
Fix for DROP TABLE/INDEX in rolled-back transaction (Hiroshi)
Fix psql crash from \l+ if MULTIBYTE enabled (Peter E)
Fix truncation of rule names during CREATE VIEW (Ross Reedstrom)
Fix PL/Perl (Alex Kapranoff)
Disallow LOCK on views (Mark Hollomon)
Disallow INSERT/UPDATE/DELETE on views (Mark Hollomon)
Disallow DROP RULE, CREATE INDEX, TRUNCATE on views (Mark Hollomon)
Allow PL/pgSQL accept non-ASCII identifiers (Tatsuo)
Allow views to proper handle GROUP BY, aggregates, DISTINCT (Tom)
Fix rare failure with TRUNCATE command (Tom)
Allow UNION/INTERSECT/EXCEPT to be used with ALL, subqueries, views,
DISTINCT, ORDER BY, SELECT...INTO (Tom)
Fix parser failures during aborted transactions (Tom)
Allow temporary relations to properly clean up indexes (Bruce)
Fix VACUUM problem with moving rows in same page (Tom)
Modify pg_dump to better handle user-defined items in template (Philip)
Allow LIMIT in VIEW (Tom)
Require cursor FETCH to honor LIMIT (Tom)
Allow PRIMARY/FOREIGN Key definitions on inherited columns (Stephan)
Allow ORDER BY, LIMIT in subqueries (Tom)
Allow UNION in CREATE RULE (Tom)
Make ALTER/DROP TABLE rollback-able (Vadim, Tom)
Store initdb collation in pg_control so collation cannot be changed (Tom)
Fix INSERT...SELECT with rules (Tom)
Fix FOR UPDATE inside views and subselects (Tom)
Fix OVERLAPS operators conform to SQL92 spec regarding NULLs (Tom)
Fix lpad() and rpad() to handle length less than input string (Tom)
Fix use of NOTIFY in some rules (Tom)
Overhaul btree code (Tom)
Fix NOT NULL use in PL/pgSQL variables (Tom)
Overhaul GIST code (Oleg)
Fix CLUSTER to preserve constraints and column default (Tom)
Improved deadlock detection handling (Tom)
Allow multiple SERIAL columns in a table (Tom)
Prevent occasional index corruption (Vadim)

Enhancements
-------------
Add OUTER JOINs (Tom)
Function manager overhaul (Tom)
Allow ALTER TABLE RENAME on indexes (Tom)
Improve CLUSTER (Tom)
Improve ps status display for more platforms (Peter E, Marc)
Improve CREATE FUNCTION failure message (Ross)
JDBC improvements (Peter, Travis Bauer, Christopher Cain, William Webber, Gunnar)
Grand Unified Configuration scheme/GUC. Many options can now be set in
data/postgresql.conf, postmaster/postgres flags, or SET commands (Peter E)
Improved handling of file descriptor cache (Tom)
New warning code about auto-created table alias entries (Bruce)
Overhaul initdb process (Tom, Peter E)
Overhaul of inherited tables; inherited tables now accessed by default;
new ONLY key word prevents it (Chris Bitmead, Tom)
ODBC cleanups/improvements (Nick Gorham, Stephan Szabo, Zoltan Kovacs, Michael Fork)
Allow renaming of temp tables (Tom)
Overhaul memory manager contexts (Tom)
pg_dumpall uses CREATE USER or CREATE GROUP rather using COPY (Peter E)
Overhaul pg_dump (Philip Warner)
Allow pg_hba.conf secondary password file to specify only username (Peter E)
Allow TEMPORARY or TEMP key word when creating temporary tables (Bruce)
New memory leak checker (Karel)
New SET SESSION CHARACTERISTICS (Thomas)
Allow nested block comments (Thomas)
Add WITHOUT TIME ZONE type qualifier (Thomas)
New ALTER TABLE ADD CONSTRAINT (Stephan)
Use NUMERIC accumulators for INTEGER aggregates (Tom)
New memory aggregate code (Tom)
New VARIANCE and STDDEV() aggregates
Improve dependency ordering of pg_dump (Philip)
New pg_restore command (Philip)
New pg_dump tar output option (Philip)
New pg_dump of large objects (Philip)
New ESCAPE option to LIKE (Thomas)
New case-insensitive LIKE - ILIKE (Thomas)
Allow functional indexes to use binary-compatible type (Tom)
Allow SQL functions to be used in more contexts (Tom)
New pg_config utility (Peter E)
New PL/pgSQL EXECUTE command which allows dynamic SQL and utility statements (Jan)
New PL/pgSQL GET DIAGNOSTICS statement for SPI value access (Jan)
New quote_identifiers() and quote_literal() functions (Jan)
New ALTER TABLE table OWNER TO user command (Mark Hollomon)
Allow subselects in FROM, i.e. FROM (SELECT ...) [AS] alias (Tom)
Update PyGreSQL to version 3.1 (D'Arcy)
Store tables as files named by OID (Vadim)
New SQL function setval(seq,val,bool) for use in pg_dump (Philip)
Require DROP VIEW to remove views, no DROP TABLE (Mark)
Allow DROP VIEW view1, view2 (Mark)
Allow multiple objects in DROP INDEX, DROP RULE, and DROP TYPE (Tom)
Allow automatic conversion to/from Unicode (Tatsuo, Eiji)
New /contrib/pgcrypto hashing functions (Marko Kreen)
New pg_dumpall --globals-only option (Peter E)
New CHECKPOINT command for WAL which creates new WAL log file (Vadim)
New AT TIME ZONE syntax (Thomas)
Allow location of Unix domain socket to be configurable (David J. MacKenzie)
Allow postmaster to listen on a specific IP address (David J. MacKenzie)
Allow socket path name to be specified in hostname by using leading slash (David J. MacKenzie)
Allow CREATE DATABASE to specify template database (Tom)
New utility to convert MySQL schema dumps to SQL92 and PostgreSQL (Thomas)
New /contrib/rserv replication toolkit (Vadim)
New file format for COPY BINARY (Tom)
New /contrib/oid2name to map numeric files to table names (B Palmer)
New "idle in transaction" ps status message (Marc)
Update to pgaccess 0.98.7 (Constantin Teodorescu)
pg_ctl now defaults to -w (wait) on shutdown, new -l (log) option
Add rudimentary dependency checking to pg_dump (Philip)

Types
-----
Fix INET/CIDR type ordering and add new functions (Tom)
Make OID behave as an unsigned type (Tom)
Allow BIGINT as synonym for INT8 (Peter E)
New int2 and int8 comparison operators (Tom)
New BIT and BIT VARYING types (Adriaan Joubert, Tom, Peter E)
CHAR() no longer faster than VARCHAR() because of TOAST (Tom)
New GIST seg/cube examples (Gene Selkov)
Improved round(numeric) handling (Tom)
Fix CIDR output formatting (Tom)
New CIDR abbrev() function (Tom)

Performance
------------
Write-Ahead Log (WAL) to provide crash recovery with less performance overhead (Vadim)
ANALYZE stage of VACUUM no longer exclusively locks table (Bruce)
Reduced file seeks (Denis Perchine)
Improve BTREE code for duplicate keys (Tom)
Store all large objects in a single table (Denis Perchine, Tom)
Improve memory allocation performance (Karel, Tom)

Source Code
---------
New function manager call conventions (Tom)
SGI portability fixes (David Kaelbling)
New configure --enable-syslog option (Peter E)
New BSDI README (Bruce)
configure script moved to top level, not /src (Peter E)
Makefile/configuration/compilation overhaul (Peter E)
New configure --with-python option (Peter E)
Solaris cleanups (Peter E)
Overhaul /contrib Makefiles (Karel)
New OpenSSL configuration option (Magnus, Peter E)
AIX fixes (Andreas)
QNX fixes (Maurizio)
New heap_open(), heap_openr() API (Tom)
Remove colon and semi-colon operators (Thomas)
New pg_class.relkind value for views (Mark Hollomon)
Rename ichar() to chr() (Karel)
New documentation for btrim(), ascii(), chr(), repeat() (Karel)
Fixes for NT/Cygwin (Pete Forman)
AIX port fixes (Andreas)
New BeOS port (David Reid, Cyril Velter)
Add proofreader's changes to docs (Addison-Wesley, Bruce)
New Alpha spinlock code (Adriaan Joubert, Compaq)
UnixWare port overhaul (Peter E)
New macOS (Darwin) port (Peter Bierman, Bruce Hartzler)
New FreeBSD Alpha port (Alfred)
Overhaul shared memory segments (Tom)
Add IBM S/390 support (Neale Ferguson)
Moved macmanuf to /contrib (Larry Rosenman)
Syslog improvements (Larry Rosenman)
New template0 database that contains no user additions (Tom)
New /contrib/cube and /contrib/seg GIST sample code (Gene Selkov)
Allow NetBSD's libedit instead of readline (Peter)
Improved assembly language source code format (Bruce)
New contrib/pg_logger
New --template option to createdb
New contrib/pg_control utility (Oliver)
New FreeBSD tools ipc_check, start-scripts/freebsd

E.347. Release 7.0.3

Release date: 2000-11-11

This has a variety of fixes from 7.0.2.

E.347.1. Migration to Version 7.0.3

A dump/restore is not required for those running 7.0.*.

E.347.2. Changes

Jdbc fixes (Peter)
Large object fix (Tom)
Fix lean in COPY WITH OIDS leak (Tom)
Fix backwards-index-scan (Tom)
Fix SELECT ... FOR UPDATE so it checks for duplicate keys (Hiroshi)
Add --enable-syslog to configure (Marc)
Fix abort transaction at backend exit in rare cases (Tom)
Fix for psql \l+ when multibyte enabled (Tatsuo)
Allow PL/pgSQL to accept non ascii identifiers (Tatsuo)
Make vacuum always flush buffers (Tom)
Fix to allow cancel while waiting for a lock (Hiroshi)
Fix for memory allocation problem in user authentication code (Tom)
Remove bogus use of int4out() (Tom)
Fixes for multiple subqueries in COALESCE or BETWEEN (Tom)
Fix for failure of triggers on heap open in certain cases (Jeroen van Vianen)
Fix for erroneous selectivity of not-equals (Tom)
Fix for erroneous use of strcmp() (Tom)
Fix for bug where storage manager accesses items beyond end of file (Tom)
Fix to include kernel errno message in all smgr elog messages (Tom)
Fix for "." not in PATH at build time (SL Baur)
Fix for out-of-file-descriptors error (Tom)
Fix to make pg_dump dump 'iscachable' flag for functions (Tom)
Fix for subselect in targetlist of Append node (Tom)
Fix for mergejoin plans (Tom)
Fix TRUNCATE failure on relations with indexes (Tom)
Avoid database-wide restart on write error (Hiroshi)
Fix nodeMaterial to honor chgParam by recomputing its output (Tom)
Fix VACUUM problem with moving chain of update row versions when source and destination of a row version lie on the same page (Tom)
Fix user.c CommandCounterIncrement (Tom)
Fix for AM/PM boundary problem in to_char() (Karel Zak)
Fix TIME aggregate handling (Tom)
Fix to_char() to avoid coredump on NULL input (Tom)
Buffer fix (Tom)
Fix for inserting/copying longer multibyte strings into char() data types (Tatsuo)
Fix for crash of backend, on abort (Tom)

E.348. Release 7.0.2

Release date: 2000-06-05

This is a repackaging of 7.0.1 with added documentation.

E.348.1. Migration to Version 7.0.2

A dump/restore is not required for those running 7.*.

E.348.2. Changes

Added documentation to tarball.
E.349. Release 7.0.1

Release date: 2000-06-01

This is a cleanup release for 7.0.

E.349.1. Migration to Version 7.0.1

A dump/restore is not required for those running 7.0.

E.349.2. Changes

Fix many CLUSTER failures (Tom)
Allow ALTER TABLE RENAME works on indexes (Tom)
Fix plpgsql to handle datetime->timestamp and timespan->interval (Bruce)
New configure --with-setproctitle switch to use setproctitle() (Marc, Bruce)
Fix the off by one errors in ResultSet from 6.5.3, and more.
jdbc ResultSet fixes (Joseph Shraibman)
optimizer tunings (Tom)
Fix create user for pgaccess
Fix for UNLISTEN failure
IRIX fixes (David Kaelbling)
QNX fixes (Andreas Kardos)
Reduce COPY IN lock level (Tom)
Change libpqeasy to use PQconnectdb() style parameters (Bruce)
Fix pg_dump to handle OID indexes (Tom)
Fix small memory leak (Tom)
Solaris fix for createdb/dropdb (Tatsuo)
Fix for non-blocking connections (Alfred Perlstein)
Fix improper recovery after RENAME TABLE failures (Tom)
Copy pg_ident.conf.sample into /lib directory in install (Bruce)
Add SJIS UDC (NEC selection IBM kanji) support (Eiji Tokuya)
Fix too long syslog message (Tatsuo)
Fix problem with quoted indexes that are too long (Tom)
JDBC ResultSet.getTimestamp() fix (Gregory Krasnow & Floyd Marinescu)
ecpg changes (Michael)

E.350. Release 7.0

Release date: 2000-05-08

This release contains improvements in many areas, demonstrating the continued growth of PostgreSQL. There are more improvements and fixes in 7.0 than in any previous release. The developers have confidence that this is the best release yet; we do our best to put out only solid releases, and this one is no exception.

Major changes in this release:

Foreign Keys

Foreign keys are now implemented, with the exception of PARTIAL MATCH foreign keys. Many users have been asking for this feature, and we are pleased to offer it.
Optimizer Overhaul

Continuing on work started a year ago, the optimizer has been improved, allowing better query plan selection and faster performance with less memory usage.

Updated psql

psql, our interactive terminal monitor, has been updated with a variety of new features. See the psql manual page for details.

Join Syntax

SQL92 join syntax is now supported, though only as INNER JOIN for this release. JOIN, NATURAL JOIN, JOIN/USING, and JOIN/ON are available, as are column correlation names.

E.350.1. Migration to Version 7.0

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release of PostgreSQL. For those upgrading from 6.5.*, you can instead use pg_upgrade to upgrade to this release; however, a full dump/reload installation is always the most robust method for upgrades.

Interface and compatibility issues to consider for the new release include:

- The date/time types datetime and timespan have been superseded by the SQL92-defined types timestamp and interval. Although there has been some effort to ease the transition by allowing PostgreSQL to recognize the deprecated type names and translate them to the new type names, this mechanism cannot be completely transparent to your existing application.

- The optimizer has been substantially improved in the area of query cost estimation. In some cases, this will result in decreased query times as the optimizer makes a better choice for the preferred plan. However, in a small number of cases, usually involving pathological distributions of data, your query times might go up. If you are dealing with large amounts of data, you might want to check your queries to verify performance.

- The JDBC and ODBC interfaces have been upgraded and extended.

- The string function CHAR_LENGTH is now a native function. Previous versions translated this into a call to LENGTH, which could result in ambiguity with other types implementing LENGTH such as the geometric types.

E.350.2. Changes

Bug Fixes

Prevent function calls exceeding maximum number of arguments (Tom)
Improve CASE construct (Tom)
Fix SELECT coalesce(f1,0) FROM int4_tbl GROUP BY f1 (Tom)
Fix SELECT sentence.words[0] FROM sentence GROUP BY sentence.words[0] (Tom)
Fix GROUP BY scan bug (Tom)
Improvements in SQL grammar processing (Tom)
Fix for views involved in INSERT ... SELECT ... (Tom)
Fix for SELECT a/2, a/2 FROM test_missing_target GROUP BY a/2 (Tom)
Fix for subselects in INSERT ... SELECT (Tom)
Prevent INSERT ... SELECT ... ORDER BY (Tom)
Fixes for relations greater than 2GB, including vacuum
Improve propagating system table changes to other backends (Tom)
Improve propagating user table changes to other backends (Tom)
Fix handling of temp tables in complex situations (Bruce, Tom)
Allow table locking at table open, improving concurrent reliability (Tom)
Properly quote sequence names in pg_dump (Ross J. Reedstrom)
Prevent DROP DATABASE while others accessing
Prevent any rows from being returned by GROUP BY if no rows processed (Tom)
Fix SELECT COUNT(1) FROM table WHERE ...' if no rows matching WHERE (Tom)
Fix pg_upgrade so it works for MVCC (Tom)
Fix for SELECT ... WHERE x IN (SELECT ... HAVING SUM(x) > 1) (Tom)
Fix problems with CURRENT_DATE used in DEFAULT (Tom)
Allow comment-only lines, and ;;; lines too. (Tom)
Improve recovery after failed disk writes, disk full (Hiroshi)
Fix cases where table is mentioned in FROM but not joined (Tom)
Allow HAVING clause without aggregate functions (Tom)
Fix for "---" comment and no trailing newline, as seen in perl interface
Improve pg_dump failure error reports (Bruce)
Allow sorts and hashes to exceed 2GB file sizes (Tom)
Fix for pg_dump dumping of inherited rules (Tom)
Fix for NULL handling comparisons (Tom)
Fix inconsistent state caused by failed CREATE/DROP commands (Hiroshi)
Fix for dbname with dash
Prevent DROP INDEX from interfering with other backends (Tom)
Fix file descriptor leak in verify_password() (Tom)
Fix for "Unable to identify an operator =\$" problem (null)
Fix ODBC so no segfault if CommLog and Debug enabled (Dirk Niggemann)
Fix for recursive exit call (Massimo)
Fix for extra-long timezones (Jeroen van Vianen)
Make pg_dump preserve primary key information (Peter E)
Prevent databases with single quotes (Peter E)
Prevent DROP DATABASE inside transaction (Peter E)
cepg memory leak fixes (Stephen Birch)
Fix for SELECT null::text, SELECT int4fac(null) and SELECT 2 + (null) (Tom)
Y2K timestamp fix (Massimo)
Fix for VACUUM 'HEAP_MOVED_IN was not expected' errors (Tom)
Fix for views with tables/columns containing spaces (Tom)
Prevent privileges on indexes (Peter E)
Fix for spinlock stuck problem when error is generated (Hiroshi)
Fix ipcclean on Linux
Fix handling of NULL constraint conditions (Tom)
Fix memory leak in odbc driver (Nick Gorham)
Fix for privilege check on UNION tables (Tom)
Fix to allow SELECT 'a' LIKE 'a' (Tom)
Fix for SELECT 1 + NULL (Tom)
Fixes to CHAR
Fix log() on numeric type (Tom)
Deprecate ':' and ':' operators
Allow vacuum of temporary tables
Disallow inherited columns with the same name as new columns
Recover or force failure when disk space is exhausted (Hiroshi)
Fix INSERT INTO ... SELECT with AS columns matching result columns
Fix INSERT ... SELECT ... GROUP BY groups by target columns not source columns (Tom)
Fix CREATE TABLE test (a char(5) DEFAULT text '', b int4) with
    INSERT (Tom)
Fix UNION with LIMIT
Fix CREATE TABLE x AS SELECT 1 UNION SELECT 2
Fix CREATE TABLE test(col char(2) DEFAULT user)
Fix mismatched types in CREATE TABLE ... DEFAULT
Fix SELECT * FROM pg_class where oid in (0,-1)
Fix SELECT COUNT('asdf') FROM pg_class WHERE oid=12
Prevent user who can create databases can modifying pg_database
    table (Peter E)
Fix btree to give a useful elog when key > 1/2 (page - overhead)
    (Tom)
Fix INSERT of 0.0 into DECIMAL(4,4) field (Tom)

Enhancements
------------
New CLI interface include file sqlcli.h, based on SQL3/SQl98
Remove all limits on query length, row length limit still exists
    (Tom)
Update jdbc protocol to 2.0 (Jens Glaser <jens@jens.de>)
Add TRUNCATE command to quickly truncate relation (Mike Mascari)
Fix to give super user and createdb user proper update catalog
    rights (Peter E)
Allow ecpg bool variables to have NULL values (Christof)
Issue ecpg error if NULL value for variable with no NULL indicator
    (Christof)
Allow ^C to cancel COPY command (Massimo)
Add SET FSYNC and SHOW PG_OPTIONS commands (Massimo)
Function name overloading for dynamically-loaded C functions
    (Frankpitt)
Add CmdTuples() to libpq++(Vince)
New CREATE CONSTRAINT TRIGGER and SET CONSTRAINTS commands (Jan)
Allow CREATE FUNCTION/WITH clause to be used for all language types
configure --enable-debug adds -g (Peter E)
configure --disable-debug removes -g (Peter E)
Allow more complex default expressions (Tom)
First real FOREIGN KEY constraint trigger functionality (Jan)
Add FOREIGN KEY ... MATCH FULL ... ON DELETE CASCADE (Jan)
Add FOREIGN KEY ... MATCH <unspecified> referential actions (Don
Baccus)
Allow WHERE restriction on ctid (physical heap location) (Hirosi)
Move pginterface from contrib to interface directory, rename to
    pgeasy (Bruce)
Change pgeasy connectdb() parameter ordering (Bruce)
Require SELECT DISTINCT target list to have all ORDER BY columns
    (Tom)
Add Oracle's COMMENT ON command (Mike Mascari <mascarim@yahoo.com>)
libpq's PQsetNoticeProcessor function now returns previous
    hook (Peter E)
Prevent PQsetNoticeProcessor from being set to NULL (Peter E)
Make USING in COPY optional (Bruce)
Allow subselects in the target list (Tom)
Allow subselects on the left side of comparison operators (Tom)
New parallel regression test (Jan)
Change backend-side COPY to write files with permissions 644 not
    666 (Tom)
Force permissions on PGDATA directory to be secure, even if it
    exists (Tom)
Release Notes

Added psql LASTOID variable to return last inserted oid (Peter E)
Allow concurrent vacuum and remove pg_vlock vacuum lock file (Tom)
Add privilege check for vacuum (Peter E)
New libpq functions to allow asynchronous connections:
PQconnectStart(),
PQconnectPoll(), PQresetStart(), PQresetPoll(), PQsetenvStart(),
PQsetenvPoll(), PQsetenvAbort (Ewan Mellor)
New libpq PQsetenv() function (Ewan Mellor)
create/alter user extension (Peter E)
New postmaster.pid and postmaster.opts under $PGDATA (Tatsuo)
New scripts for create/drop user/db (Peter E)
Major psql overhaul (Peter E)
Add const to libpq interface (Peter E)
New libpq function PQoidValue (Peter E)
Show specific non-aggregate causing problem with GROUP BY (Tom)
Make changes to pg_shadow recreate pg_pwd file (Peter E)
Add aggregate(DISTINCT ...) (Tom)
Allow flag to control COPY input/output of NULLs (Peter E)
Make postgres user have a password by default (Peter E)
Add CREATE/ALTER/DROP GROUP (Peter E)
All administration scripts now support --long options (Peter E, Karel)
Vacuumdb script now supports --all option (Peter E)
cpq new portable FETCH syntax
Add ecpg EXEC SQL IFDEF, EXEC SQL IFNDEF, EXEC SQL ELSE, EXEC SQL ELIF
and EXEC SQL ENDIF directives
Add pg_ctl script to control backend start-up (Tatsuo)
Add postmaster.opts.default file to store start-up flags (Tatsuo)
Allow --with-mb=SQL_ASCII
Increase maximum number of index keys to 16 (Bruce)
Increase maximum number of function arguments to 16 (Bruce)
Allow configuration of maximum number of index keys and arguments (Bruce)
Allow unprivileged users to change their passwords (Peter E)
Password authentication enabled; required for new users (Peter E)
Disallow dropping a user who owns a database (Peter E)
Change initdb option --with-mb to --enable-multibyte
Add option for initdb to prompts for superuser password (Peter E)
Allow complex type casts like col::numeric(9,2) and col::int2::float8 (Tom)
Updated user interfaces on initdb, initlocation, pg_dump, ipcclean (Peter E)
New pg_char_to_encoding() and pg_encoding_to_char() functions (Tatsuo)
libpq non-blocking mode (Alfred Perlstein)
Improve conversion of types in casts that don't specify a length
New plperl internal programming language (Mark Hollomon)
Allow COPY IN to read file that do not end with a newline (Tom)
Indicate when long identifiers are truncated (Tom)
Allow aggregates to use type equivalency (Peter E)
Add Oracle's to_char(), to_date(), to_datetime(), to_timestamp(), to_number()
conversion functions (Karel Zak <zakkr@zf.jcu.cz>)
Add SELECT DISTINCT ON (expr [, expr ...]) targetlist ... (Tom)
Check to be sure ORDER BY is compatible with the DISTINCT operation (Tom)
Add NUMERIC and int8 types to ODBC
Release Notes

Improve EXPLAIN results for Append, Group, Agg, Unique (Tom)
Add ALTER TABLE ... ADD FOREIGN KEY (Stephan Szabo)
Allow SELECT .. FOR UPDATE in PL/pgSQL (Hirosi)
Enable backward sequential scan even after reaching EOF (Hirosi)
Add btree indexing of boolean values, >= and <= (Don Baccus)
Print current line number when COPY FROM fails (Massimo)
Recognize POSIX time zone e.g. "PST+8" and "GMT-8" (Thomas)
Add DEC as synonym for DECIMAL (Thomas)
Add SESSION_USER as SQL92 key word, same as CURRENT_USER (Thomas)
Implement SQL92 column aliases (aka correlation names) (Thomas)
Implement SQL92 join syntax (Thomas)
Make INTERVAL reserved word allowed as a column identifier (Thomas)
Implement REINDEX command (Hirosi)
Accept ALL in aggregate function SUM(ALL col) (Tom)
Prevent GROUP BY from using column aliases (Tom)
New psql \encoding option (Tatsuo)
Allow PQrequestCancel() to terminate when in waiting-for-lock state (Hirosi)
Allow negation of a negative number in all cases
Add ecpg descriptors (Christof, Michael)
Allow CREATE VIEW v AS SELECT f1::char(8) FROM tbl
Allow casts with length, like foo::char(8)
New libpq functions PQsetClientEncoding(), PQclientEncoding()
(Tatsuo)
Add support for SJIS user defined characters (Tatsuo)
Larger views/rules supported
Make libpq's PQconndefaults() thread-safe (Tom)
Disable // as comment to be ANSI conforming, should use -- (Tom)
Allow column aliases on views CREATE VIEW name (collist)
Fixes for views with subqueries (Tom)
Allow UPDATE table SET fld = (SELECT ...) (Tom)
SET command options no longer require quotes
Update pgaccess to 0.98.6
New SET SEED command
New pg_options.sample file
New SET FSYNC command (Massimo)
Allow pg_descriptions when creating tables
Allow pg_descriptions when creating types, columns, and functions
Allow psql \copy to allow delimiters (Peter E)
Allow psql to print nulls as distinct from "" [null] (Peter E)

Types
-----
Many array fixes (Tom)
Allow bare column names to be subscripted as arrays (Tom)
Improve type casting of int and float constants (Tom)
Cleanup for int8 inputs, range checking, and type conversion (Tom)
Fix for SELECT timespan('21:11:26':::time) (Tom)
netmask('x.x.x.x/0') is 255.255.255.255 instead of 0.0.0.0 (Oleg Sharoiko)
Add btree index on NUMERIC (Jan)
Perl fix for large objects containing NUL characters (Douglas Thomson)
ODBC fix for large objects (free)
Fix indexing of cidr data type
Fix for Ethernet MAC addresses (macaddr type) comparisons
Fix for date/time types when overflows happened in computations (Tom)
Release Notes

Allow array on int8 (Peter E)
Fix for rounding/overflow of NUMERIC type, like NUMERIC(4,4) (Tom)
Allow NUMERIC arrays
Fix bugs in NUMERIC ceil() and floor() functions (Tom)
Make char_length()/octet_length including trailing blanks (Tom)
Made abstime/reltime use int4 instead of time_t (Peter E)
New ltext data type for compressed text fields
Revise code to handle coercion of int and float constants (Tom)
Start at new code to implement a BIT and BIT VARYING type (Adriaan Joubert)
NUMERIC now accepts scientific notation (Tom)
NUMERIC to int4 rounds (Tom)
Convert float4/8 to NUMERIC properly (Tom)
Allow type conversion with NUMERIC (Thomas)
Make ISO date style (2000-02-16 09:33) the default (Thomas)
Add NATIONAL CHAR [ VARYING ] (Thomas)
Allow NUMERIC round and trunc to accept negative scales (Tom)
New TIME WITH TIME ZONE type (Thomas)
Add MAX()/MIN() on time type (Thomas)
Add abs(), mod(), fac() for int8 (Thomas)
Rename functions to round(), sqrt(), cbrt(), pow() for float8 (Thomas)
Add transcendental math functions (e.g. sin(), acos()) for float8 (Thomas)
Add exp() and ln() for NUMERIC type
Rename NUMERIC power() to pow() (Thomas)
Improved TRANSLATE() function (Edwin Ramirez, Tom)
Allow X=-Y operators (Tom)
Allow SELECT float8(COUNT(*))/(SELECT COUNT(*) FROM t) FROM t GROUP BY f1; (Tom)
Allow LOCALE to use indexes in regular expression searches (Tom)
Allow creation of functional indexes to use default types

Performance
-----------
Prevent exponential space consumption with many AND's and OR's (Tom)
Collect attribute selectivity values for system columns (Tom)
Reduce memory usage of aggregates (Tom)
Fix for LIKE optimization to use indexes with multibyte encodings (Tom)
Fix r-tree index optimizer selectivity (Thomas)
Improve optimizer selectivity computations and functions (Tom)
Optimize btree searching for cases where many equal keys exist (Tom)
Enable fast LIKE index processing only if index present (Tom)
Re-use free space on index pages with duplicates (Tom)
Improve hash join processing (Tom)
Prevent descending sort if result is already sorted (Hiroshi)
Allow commuting of index scan query qualifications (Tom)
Prefer index scans in cases where ORDER BY/GROUP BY is required (Tom)
Allocate large memory requests in fix-sized chunks for performance (Tom)
Fix vacuum's performance by reducing memory allocation requests (Tom)
Implement constant-expression simplification (Bernard Frankpitt, Tom)
Use secondary columns to be used to determine start of index scan (Hiroshi)
Prevent quadruple use of disk space when doing internal sorting (Tom)
Faster sorting by calling fewer functions (Tom)
Create system indexes to match all system caches (Bruce, Hiroshi)
Make system caches use system indexes (Bruce)
Make all system indexes unique (Bruce)
Improve pg_statistic management for VACUUM speed improvement (Tom)
Flush backend cache less frequently (Tom, Hiroshi)
COPY now reuses previous memory allocation, improving performance (Tom)
Improve optimization cost estimation (Tom)
Improve optimizer estimate of range queries \( x > \) lowbound AND \( x < \) highbound (Tom)
Use DNF instead of CNF where appropriate (Tom, Taral)
Further cleanup for OR-of-AND WHERE-clauses (Tom)
Make use of index in OR clauses \( (x = 1 \text{ AND } y = 2) \text{ OR } (x = 2 \text{ AND } y = 4) \) (Tom)
Smarter optimizer computations for random index page access (Tom)
New SET variable to control optimizer costs (Tom)
Optimizer queries based on LIMIT, OFFSET, and EXISTS qualifications (Tom)
Reduce optimizer internal housekeeping of join paths for speedup (Tom)
Major subquery speedup (Tom)
Fewer fsync writes when fsync is not disabled (Tom)
Improved LIKE optimizer estimates (Tom)
Prevent fsync in SELECT-only queries (Vadim)
Make index creation use psort code, because it is now faster (Tom)
Allow creation of sort temp tables > 1 Gig

Source Tree Changes
-------------------
Fix for linux PPC compile
New generic expression-tree-walker subroutine (Tom)
Change form() to varargform() to prevent portability problems
Improved range checking for large integers on Alphas
Clean up #include in /include directory (Bruce)
Add scripts for checking includes (Bruce)
Remove un-needed #include's from *.c files (Bruce)
Change #include's to use <> and "" as appropriate (Bruce)
Enable Windows compilation of libpq
Alpha spinlock fix from Uncle George <gatgul@voicenet.com>
Overhaul of optimizer data structures (Tom)
Fix to cygipc library (Yutaka Tanida)
Allow pgsql to work on newer Cygwin snapshots (Dan)
New catalog version number (Tom)
Add Linux ARM
Rename heap_replace to heap_update
Update for QNX (Dr. Andreas Kardos)
New platform-specific regression handling (Tom)
Rename oid8 -> oidvector and int28 -> int2vector (Bruce)
Included all yacc and lex files into the distribution (Peter E.)
Remove lextest, no longer needed (Peter E)
Fix for libpq and psql on Windows (Magnus)
Internally change datetime and timespan into timestamp and interval (Thomas)
Fix for plpgsql on BSD/OS
Add SQL_ASCII test case to the regression test (Tatsuo)
configure --with-mb now deprecated (Tatsuo)
NT fixes
NetBSD fixes (Johnny C. Lam <lamj@stat.cmu.edu>)
Fixes for Alpha compiles
New multibyte encodings

E.351. Release 6.5.3

Release date: 1999-10-13

This is basically a cleanup release for 6.5.2. We have added a new PgAccess that was missing in 6.5.2, and installed an NT-specific fix.

E.351.1. Migration to Version 6.5.3

A dump/restore is not required for those running 6.5.*.

E.351.2. Changes

Updated version of pgaccess 0.98
NT-specific patch
Fix dumping rules on inherited tables

E.352. Release 6.5.2

Release date: 1999-09-15

This is basically a cleanup release for 6.5.1. We have fixed a variety of problems reported by 6.5.1 users.

E.352.1. Migration to Version 6.5.2

A dump/restore is not required for those running 6.5.*.

E.352.2. Changes

subselect+CASE fixes(Tom)
Add SHLIB_LINK setting for solaris_i386 and solaris_sparc ports(Daren Sefcik)
Fixes for CASE in WHERE join clauses(Tom)
Fix BTScan abort(Tom)
Repair the check for redundant UNIQUE and PRIMARY KEY indexes(Thomas)
Improve it so that it checks for multicoloumn constraints(Thomas)
Fix for Windows making problem with MB enabled(Hiroki Kataoka)
Allow BSD yacc and bison to compile pl code(Bruce)
Fix SET NAMES working
int8 fixes(Thomas)
Fix vacuum's memory consumption(Hiroshi,Tatsuo)
Reduce the total memory consumption of vacuum(Tom)
Fix for timestamp (datetime)
Rule deparsing bugfixes (Tom)
Fix quoting problems in mkMakefile.tcldefs.sh.in and
mkMakefile.tkdefs.sh.in (Tom)
This is to re-use space on index pages freed by vacuum (Vadim)
document -x for pg_dump (Bruce)
Fix for unary operators in rule deparser (Tom)
Comment out FileUnlink of excess segments during mdtruncate() (Tom)
IRIX linking fix from Yu Cao >yucao@falcon.kla-tencor.com<
Repair logic error in LIKE: should not return LIKE_ABORT
when reach end of pattern before end of text (Tom)
Repair incorrect cleanup of heap memory allocation during
transaction abort (Tom)
Updated version of pgaccess 0.98

E.353. Release 6.5.1

Release date: 1999-07-15

This is basically a cleanup release for 6.5. We have fixed a variety of problems reported by 6.5 users.

E.353.1. Migration to Version 6.5.1

A dump/restore is not required for those running 6.5.

E.353.2. Changes

Add NT README file
Portability fixes for linux_ppc, IRIX, linux_alpha, OpenBSD, alpha
Remove QUERY_LIMIT, use SELECT...LIMIT
Fix for EXPLAIN on inheritance (Tom)
Patch to allow vacuum on multisegment tables (Hiroshi)
R-Tree optimizer selectivity fix (Tom)
ACL file descriptor leak fix (Atsushi Ogawa)
New expression subtree code (Tom)
Avoid disk writes for read-only transactions (Vadim)
Fix for removal of temp tables if last transaction was aborted (Bruce)
Fix to prevent too large row from being created (Bruce)
plpgsql fixes
Allow port numbers 32k - 64k (Bruce)
Add ^ precedence (Bruce)
Rename sort files called pg_temp to pg_sorttemp (Bruce)
Fix for microseconds in time values (Tom)
Tutorial source cleanup
New linux_m68k port
Fix for sorting of NULL's in some cases (Tom)
Shared library dependencies fixed (Tom)
Fixed glitches affecting GROUP BY in subselects (Tom)
Fix some compiler warnings (Tomoaki Nishiyama)
Add Win1250 (Czech) support (Pavel Behal)

E.354. Release 6.5

Release date: 1999-06-09
This release marks a major step in the development team's mastery of the source code we inherited from Berkeley. You will see we are now easily adding major features, thanks to the increasing size and experience of our world-wide development team.

Here is a brief summary of the more notable changes:

Multiversion concurrency control (MVCC)

This removes our old table-level locking, and replaces it with a locking system that is superior to most commercial database systems. In a traditional system, each row that is modified is locked until committed, preventing reads by other users. MVCC uses the natural multiversion nature of PostgreSQL to allow readers to continue reading consistent data during writer activity. Writers continue to use the compact pg_log transaction system. This is all performed without having to allocate a lock for every row like traditional database systems. So, basically, we no longer are restricted by simple table-level locking; we have something better than row-level locking.

Hot backups from pg_dump

pg_dump takes advantage of the new MVCC features to give a consistent database dump/backup while the database stays online and available for queries.

Numeric data type

We now have a true numeric data type, with user-specified precision.

Temporary tables

Temporary tables are guaranteed to have unique names within a database session, and are destroyed on session exit.

New SQL features

We now have CASE, INTERSECT, and EXCEPT statement support. We have new LIMIT/OFFSET, SET TRANSACTION ISOLATION LEVEL, SELECT ... FOR UPDATE, and an improved LOCK TABLE command.

Speedups

We continue to speed up PostgreSQL, thanks to the variety of talents within our team. We have sped up memory allocation, optimization, table joins, and row transfer routines.

Ports

We continue to expand our port list, this time including Windows NT/ix86 and NetBSD/arm32.

Interfaces

Most interfaces have new versions, and existing functionality has been improved.

Documentation

New and updated material is present throughout the documentation. New FAQs have been contributed for SGI and AIX platforms. The Tutorial has introductory information on SQL from Stefan Simkovics. For the User's Guide, there are reference pages covering the postmaster and more utility programs, and a new appendix contains details on date/time behavior. The Administrator's Guide has a new chapter on troubleshooting from Tom Lane. And the Programmer's Guide has a description of query processing, also from Stefan, and details on obtaining the PostgreSQL source tree via anonymous CVS and CVSup.

E.354.1. Migration to Version 6.5

A dump/restore using pg_dump is required for those wishing to migrate data from any previous release of PostgreSQL. pg_upgrade can not be used to upgrade to this release because the on-disk structure of the tables has changed compared to previous releases.
The new Multiversion Concurrency Control (MVCC) features can give somewhat different behaviors in multiuser environments. Read and understand the following section to ensure that your existing applications will give you the behavior you need.

E.354.1.1. Multiversion Concurrency Control

Because readers in 6.5 don’t lock data, regardless of transaction isolation level, data read by one transaction can be overwritten by another. In other words, if a row is returned by SELECT it doesn’t mean that this row really exists at the time it is returned (i.e. sometime after the statement or transaction began) nor that the row is protected from being deleted or updated by concurrent transactions before the current transaction does a commit or rollback.

To ensure the actual existence of a row and protect it against concurrent updates one must use SELECT FOR UPDATE or an appropriate LOCK TABLE statement. This should be taken into account when porting applications from previous releases of PostgreSQL and other environments.

Keep the above in mind if you are using contrib/refint.* triggers for referential integrity. Additional techniques are required now. One way is to use LOCK parent_table IN SHARE ROW EXCLUSIVE MODE command if a transaction is going to update/delete a primary key and use LOCK parent_table IN SHARE MODE command if a transaction is going to update/insert a foreign key.

Note

Note that if you run a transaction in SERIALizable mode then you must execute the LOCK commands above before execution of any DML statement (SELECT/INSERT/DELETE/UPDATE/FETCH/COPY_TO) in the transaction.

These inconveniences will disappear in the future when the ability to read dirty (uncommitted) data (regardless of isolation level) and true referential integrity will be implemented.

E.354.2. Changes

Bug Fixes

---

Fix text<->float8 and text<->float4 conversion functions (Thomas)
Fix for creating tables with mixed-case constraints (Billy)
Change exp()/pow() behavior to generate error on underflow/overflow (Jan)
Fix bug in pg_dump -z
Memory overrun cleanups (Tatsuo)
Fix for lo_import crash (Tatsuo)
Adjust handling of data type names to suppress double quotes (Thomas)
Use type coercion for matching columns and DEFAULT (Thomas)
Fix deadlock so it only checks once after one second of sleep (Bruce)
Fixes for aggregates and PL/pgSQL (Hiroshi)
Fix for subquery crash (Vadim)
Fix for libpq function PQfnumber and case-insensitive names (Bahman Rafatjoo)
Fix for large object write-in-middle, no extra block, memory consumption (Tatsuo)
Fix for pg_dump -d or -D and quote special characters in INSERT
Repair serious problems with dynahash (Tom)
Fix INET/CIDR portability problems
Fix problem with selectivity error in ALTER TABLE ADD COLUMN (Bruce)
Fix executor so mergejoin of different column types works (Tom)
Fix for Alpha OR selectivity bug
Fix OR index selectivity problem (Bruce)
Fix so \d shows proper length for char()/varchar() (Ryan)
Fix tutorial code (Clark)
Improve destroyuser checking (Oliver)
Fix for Kerberos (Rodney McDuff)
Fix for dropping database while dirty buffers (Bruce)
Fix so sequence nextval() can be case-sensitive (Bruce)
Fix != operator
Drop buffers before destroying database files (Bruce)
Fix case where executor evaluates functions twice (Tatsuo)
Allow sequence nextval actions to be case-sensitive (Bruce)
Fix optimizer indexing not working for negative numbers (Bruce)
Fix for memory leak in executor with fjiIsNull
Fix for aggregate memory leaks (Erik Riedel)
Allow user name containing a dash to grant privileges
Cleanup of NULL in inet types
Clean up system table bugs (Tom)
Fix problems of PAGER and \? command (Masaaki Sakaida)
Reduce default multisegment file size limit to 1GB (Peter)
Fix for dumping of CREATE OPERATOR (Tom)
Fix for backward scanning of cursors (Hiroshi Inoue)
Fix for COPY FROM STDIN when using \l (Tom)
Fix for subselect is compared inside an expression (Jan)
Fix handling of error reporting while returning rows (Tom)
Fix problems with reference to array types (Tom, Jan)
Prevent UPDATE SET oid (Jan)
Fix pg_dump so -t option can handle case-sensitive tablenames
Fixes for GROUP BY in special cases (Tom, Jan)
Fix for memory leak in failed queries (Tom)
DEFAULT now supports mixed-case identifiers (Tom)
Fix for multisegment uses of DROP/RENAME table, indexes (Ole Gjerde)
Disable use of pg_dump with both -o and -d options (Bruce)
Allow pg_dump to properly dump group privileges (Bruce)
Fix GROUP BY in INSERT INTO table SELECT * FROM table2 (Jan)
Fix for computations in views (Jan)
Fix for aggregates on array indexes (Tom)
Fix for DEFAULT handles single quotes in value requiring too many quotes
Fix security problem with non-super users importing/exporting large objects (Tom)
Rollback of transaction that creates table cleaned up properly (Tom)
Fix to allow long table and column names to generate proper serial names (Tom)

Enhancements
-----------
Add "vacuumdb" utility
Speed up libpq by allocating memory better (Tom)
EXPLAIN all indexes used (Tom)
Implement CASE, COALESCE, NULLIF expression (Thomas)
New pg_dump table output format (Constantin)
Add string min()/max() functions (Thomas)
Extend new type coercion techniques to aggregates (Thomas)
New moddatetime contrib (Terry)
Update to pgaccess 0.96 (Constantin)
Add routines for single-byte "char" type (Thomas)
Improved substr() function (Thomas)
Improved multibyte handling (Tatsuo)
Multiversion concurrency control/MVCC (Vadim)
New Serialized mode (Vadim)
Fix for tables over 2gigs (Peter)
New SET TRANSACTION ISOLATION LEVEL (Vadim)
New LOCK TABLE IN ... MODE (Vadim)
Update ODBC driver (Byron)
New NUMERIC data type (Jan)
New SELECT FOR UPDATE (Vadim)
Handle "NaN" and "Infinity" for input values (Jan)
Improved date/year handling (Thomas)
Improved handling of backend connections (Magnus)
New options ELOG_TIMESTAMP and USE_SYSLOG options for log files (Massimo)
New TCL_ARRAYS option (Massimo)
New INTERSECT and EXCEPT (Stefan)
New pg_index.indisprimary for primary key tracking (D'Arcy)
New pg_dump option to allow dropping of tables before creation (Brook)
Speedup of row output routines (Tom)
New READ COMMITTED isolation level (Vadim)
New TEMP tables/indexes (Bruce)
Prevent sorting if result is already sorted (Jan)
New memory allocation optimization (Jan)
Allow psql to do \p\g (Bruce)
Allow multiple rule actions (Jan)
Added LIMIT/OFFSET functionality (Jan)
Improve optimizer when joining a large number of tables (Bruce)
New intro to SQL from S. Simkovics' Master's Thesis (Stefan, Thomas)
New intro to backend processing from S. Simkovics' Master's Thesis (Stefan)
Improved int8 support (Ryan Bradetich, Thomas, Tom)
New routines to convert between int8 and text/varchar types (Thomas)
New bushy plans, where meta-tables are joined (Bruce)
Enable right-hand queries by default (Bruce)
Allow reliable maximum number of backends to be set at configure time
(--with-maxbackends and postmaster switch (--N backends)) (Tom)
GEQO default now 10 tables because of optimizer speedups (Tom)
Allow NULL=Var for MS-SQL portability (Michael, Bruce)
Modify contrib check_primary_key() so either "automatic" or "dependent" (Anand)
Allow psql \d on a view show query (Ryan)
Speedup for LIKE (Bruce)
Ecpg fixes/features, see src/interfaces/ecpg/ChangeLog file (Michael)
JDBC fixes/features, see src/interfaces/jdbc/CHANGECLOG (Peter)
Make % operator have precedence like / (Bruce)
Add new postgres -O option to allow system table structure changes (Bruce)
Update contrib/pginterface/findoidjoins script (Tom)
Major speedup in vacuum of deleted rows with indexes (Vadim)
Allow non-SQL functions to run different versions based on arguments (Tom)
Add -E option that shows actual queries sent by \dt and friends (Masaaki Sakaida)
Add version number in start-up banners for psql (Masaaki Sakaida)
New contrib/vacuumlo removes large objects not referenced (Peter)
New initialization for table sizes so non-vacuumed tables perform better (Tom)
Improve error messages when a connection is rejected (Tom)
Support for arrays of char() and varchar() fields (Massimo)
Overhaul of hash code to increase reliability and performance (Tom)
Update to PyGreSQL 2.4 (D'Arcy)
Changed debug options so -d4 and -d5 produce different node displays (Jan)
New pg_options: pretty_plan, pretty_parse, pretty_rewritten (Jan)
Better optimization statistics for system table access (Tom)
Better handling of non-default block sizes (Massimo)
Improve GEQO optimizer memory consumption (Tom)
UNION now supports ORDER BY of columns not in target list (Jan)
Major libpq++ improvements (Vince Vielhaber)
pg_dump now uses -z (ACL's) as default (Bruce)
backend cache, memory speedups (Tom)
have pg_dump do everything in one snapshot transaction (Vadim)
fix for large object memory leakage, fix for pg_dumping (Tom)
INET type now respects netmask for comparisons
Make VACUUM ANALYZE only use a readlock (Vadim)
Allow VIEWS on UNIONS (Jan)
pg_dump now can generate consistent snapshots on active databases (Vadim)

Source Tree Changes
-------------------
Improve port matching (Tom)
Portability fixes for SunOS
Add Windows NT backend port and enable dynamic loading (Magnus and Daniel Horak)
New port to Cobalt Qube (Mips) running Linux (Tatsuo)
Port to NetBSD/m68k (Mr. Mutsuki Nakajima)
Port to NetBSD/sun3 (Mr. Mutsuki Nakajima)
Port to NetBSD/macppc (Toshimi Aoki)
Fix for tcl/tk configuration (Vince)
Removed CURRENT key word for rule queries (Jan)
NT dynamic loading now works (Daniel Horak)
Add ARM32 support (Andrew McMurry)
Better support for HP-UX 11 and UnixWare
Improve file handling to be more uniform, prevent file descriptor leak (Tom)
New install commands for plpgsql (Jan)

**E.355. Release 6.4.2**

**Release date:** 1998-12-20

The 6.4.1 release was improperly packaged. This also has one additional bug fix.

**E.355.1. Migration to Version 6.4.2**

A dump/restore is *not* required for those running 6.4.*.
E.355.2. Changes

Fix for datetime constant problem on some platforms (Thomas)

E.356. Release 6.4.1

Release date: 1998-12-18

This is basically a cleanup release for 6.4. We have fixed a variety of problems reported by 6.4 users.

E.356.1. Migration to Version 6.4.1

A dump/restore is not required for those running 6.4.

E.356.2. Changes

Add pg_dump -N flag to force double quotes around identifiers.
This is
the default (Thomas)
Fix for NOT in where clause causing crash (Bruce)
EXPLAIN VERBOSE coredump fix (Vadim)
Fix shared-library problems on Linux
Fix test for table existence to allow mixed-case and whitespace in
the table name (Thomas)
Fix a couple of pg_dump bugs
Configure matches template/.similar entries better (Tom)
Change builtin function names from SPI_* to spi_*
OR WHERE clause fix (Vadim)
Fixes for mixed-case table names (Billy)
contrib/linux/postgres.init.csh/sh fix (Thomas)
libpq memory overrun fix
SunOS fixes (Tom)
Change exp() behavior to generate error on underflow (Thomas)
pg_dump fixes for memory leak, inheritance constraints, layout change
update pgaccess to 0.93
Fix prototype for 64-bit platforms
Multibyte fixes (Tatsuo)
New ecpg man page
Fix memory overruns (Tatsuo)
Fix for lo_import() crash (Bruce)
Better search for install program (Tom)
Timezone fixes (Tom)
HP-UX fixes (Tom)
Use implicit type coercion for matching DEFAULT values (Thomas)
Add routines to help with single-byte (internal) character type (Thomas)
Compilation of libpq for Windows fixes (Magnus)
Upgrade to PyGreSQL 2.2 (D'Arcy)

E.357. Release 6.4

Release date: 1998-10-30
There are many new features and improvements in this release. Thanks to our developers and maintainers, nearly every aspect of the system has received some attention since the previous release. Here is a brief, incomplete summary:

- Views and rules are now functional thanks to extensive new code in the rewrite rules system from Jan Wieck. He also wrote a chapter on it for the *Programmer's Guide*.

- Jan also contributed a second procedural language, PL/pgSQL, to go with the original PL/pgTCL procedural language he contributed last release.

- We have optional multiple-byte character set support from Tatsuo Ishii to complement our existing locale support.

- Client/server communications has been cleaned up, with better support for asynchronous messages and interrupts thanks to Tom Lane.

- The parser will now perform automatic type coercion to match arguments to available operators and functions, and to match columns and expressions with target columns. This uses a generic mechanism which supports the type extensibility features of PostgreSQL. There is a new chapter in the *User's Guide* which covers this topic.

- Three new data types have been added. Two types, *inet* and *cidr*, support various forms of IP network, subnet, and machine addressing. There is now an 8-byte integer type available on some platforms. See the chapter on data types in the *User's Guide* for details. A fourth type, *serial*, is now supported by the parser as an amalgam of the int4 type, a sequence, and a unique index.

- Several more SQL92-compatible syntax features have been added, including INSERT DEFAULT VALUES

- The automatic configuration and installation system has received some attention, and should be more robust for more platforms than it has ever been.

**E.357.1. Migration to Version 6.4**

A dump/restore using pg_dump or pg_dumpall is required for those wishing to migrate data from any previous release of PostgreSQL.

**E.357.2. Changes**

**Bug Fixes**

----------

- Fix for a tiny memory leak in PQsetdb/PQfinish (Bryan)
- Remove char2-16 data types, use char/varchar (Darren)
- PQfn not handles a NOTICE message (Anders)
- Reduced busywaiting overhead for spinlocks with many backends (dg)
- Stuck spinlock detection (dg)
- Fix up "ISO-style" timespan decoding and encoding (Thomas)
- Fix problem with table drop after rollback of transaction (Vadim)
- Change error message and remove non-functional update message (Vadim)
- Fix for COPY array checking
- Fix for SELECT 1 UNION SELECT NULL
- Fix for buffer leaks in large object calls (Pascal)
- Change owner from oid to int4 type (Bruce)
- Fix a bug in the oracle compatibility functions btrim() ltrim() and rtrim()
- Fix for shared invalidation cache overflow (Massimo)
- Prevent file descriptor leaks in failed COPY's (Bruce)
- Fix memory leak in libpgtcl's pg_select (Constantin)
Fix problems with username/passwords over 8 characters (Tom)
Fix problems with handling of asynchronous NOTIFY in backend (Tom)
Fix of many bad system table entries (Tom)

Enhancements
-------------
Upgrade ecpg and ecpglib, see src/interfaces/ecpc/ChangeLog (Michael)
Show the index used in an EXPLAIN (Zeugswetter)
EXPLAIN invokes rule system and shows plan(s) for rewritten queries (Jan)
Multibyte awareness of many data types and functions, via configure (Tatsuo)
New configure --with-mb option (Tatsuo)
New initdb --pgencoding option (Tatsuo)
New createdb -E multibyte option (Tatsuo)
Select version(); now returns PostgreSQL version (Jeroen)
libpq now allows asynchronous clients (Tom)
Allow cancel from client of backend query (Tom)
psql now cancels query with Control-C (Tom)
libpq users need not issue dummy queries to get NOTIFY messages (Tom)
NOTIFY now sends sender's PID, so you can tell whether it was your own (Tom)
Pnresult struct now includes associated error message, if any (Tom)
Define "tz_hour" and "tz_minute" arguments to date_part () (Thomas)
Add routines to convert between varchar and bpchar (Thomas)
Add routines to allow sizing of varchar and bpchar into target columns (Thomas)
Add bit flags to support timezonehour and minute in data retrieval (Thomas)
Allow more variations on valid floating point numbers (e.g. ".1", "1e6") (Thomas)
Fixes for unary minus parsing with leading spaces (Thomas)
Implement TIMEZONE_HOUR, TIMEZONE_MINUTE per SQL92 specs (Thomas)
Check for and properly ignore FOREIGN KEY column constraints (Thomas)
Define USER as synonym for CURRENT_USER per SQL92 specs (Thomas)
Enable HAVING clause but no fixes elsewhere yet.
Make "char" type a synonym for "char(1)" (actually implemented as bpchar) (Thomas)
Save string type if specified for DEFAULT clause handling (Thomas)
Coerce operations involving different data types (Thomas)
Allow some index use for columns of different types (Thomas)
Add capabilities for automatic type conversion (Thomas)
Cleanups for large objects, so file is truncated on open (Peter)
Readline cleanups (Tom)
Allow psql \f \ to make spaces as delimiter (Bruce)
Pass pg_attribute.atttypmod to the frontend for column field lengths (Tom, Bruce)
Msgsql compatibility library in /contrib (Aldrin)
Remove the requirement that ORDER/GROUP BY clause identifiers be included in the target list (David)
Convert columns to match columns in UNION clauses (Thomas)
Remove fork () / exec () and only do fork () (Bruce)
Jdbc cleanups (Peter)
Show backend status on ps command line (only works on some platforms) (Bruce)
Pg_hba.conf now has a sameuser option in the database field
Make lo_unlink take oid param, not int4
New DISABLE_COMPLEX_MACRO for compilers that cannot handle our
macros (Bruce)
Libpgtcl now handles NOTIFY as a Tcl event, need not send dummy
queries (Tom)
libpgtcl cleanups (Tom)
Add -error option to libpgtcl's pg_result command (Tom)
New locale patch, see docs/README/locale (Oleg)
Fix for pg_dump so CONSTRAINT and CHECK syntax is correct (ccb)
New contrib/io code for large object orphan removal (Peter)
New psql command "SET CLIENT_ENCODING TO 'encoding'" for multibytes
feature, see /doc/README.mb (Tatsuo)
contrib/noupdate code to revoke update permission on a column
libpq can now be compiled on Windows (Magnus)
Add PQsetdbLogin() in libpq
New 8-byte integer type, checked by configure for OS
support (Thomas)
Better support for quoted table/column names (Thomas)
Surround table and column names with double-quotes in
pg_dump (Thomas)
PQreset() now works with passwords (Tom)
Handle case of GROUP BY target list column number out of
range (David)
Allow UNION in subselects
Add auto-size to screen to \d? commands (Bruce)
Use UNION to show all \d? results in one query (Bruce)
Add \d? field search feature (Bruce)
Pg_dump issues fewer \connect requests (Tom)
Make pg_dump -z flag work better, document it in manual page (Tom)
Add HAVING clause with full support for subselects and
unions (Stephan)
Full text indexing routines in contrib/fulltextindex (Maarten)
Transaction ids now stored in shared memory (Vadim)
New PGCLIENTENCODING when issuing COPY command (Tatsuo)
Support for SQL92 syntax "SET NAMES" (Tatsuo)
Support for LATIN2-5 (Tatsuo)
Add UNICODE regression test case (Tatsuo)
Lock manager cleanup, new locking modes for LLL (Vadim)
Allow index use with OR clauses (Bruce)
Allows "SELECT NULL ORDER BY 1;"
Explain VERBOSE prints the plan, and now pretty-prints the plan to
the postmaster log file (Bruce)
Add indexes display to \d command (Bruce)
Allow GROUP BY on functions (David)
New pg_class.relkind for large objects (Bruce)
New way to send libpq NOTICE messages to a different location (Tom)
New \w write command to psql (Bruce)
New /contrib/findoidjoins scans oid columns to find join
relationships (Bruce)
Allow binary-compatible indexes to be considered when checking for
valid
Indexes for restriction clauses containing a constant (Thomas)
New ISBN/ISSN code in /contrib/isbn_issn
Allow NOT LIKE, IN, NOT IN, BETWEEN, and NOT BETWEEN
constraint (Thomas)
New rewrite system fixes many problems with rules and views (Jan)
  * Rules on relations work
  * Event qualifications on insert/update/delete work
* New OLD variable to reference CURRENT; CURRENT will be remove in future
* Update rules can reference NEW and OLD in rule qualifications/actions
  * Insert/update/delete rules on views work
  * Multiple rule actions are now supported, surrounded by parentheses
* Regular users can create views/rules on tables they have RULE permits
  * Rules and views inherit the privileges of the creator
  * No rules at the column level
  * No UPDATE NEW/OLD rules
  * New pg_tables, pg_indexes, pg_rules and pg_views system views
  * Only a single action on SELECT rules
  * Total rewrite overhaul, perhaps for 6.5
* handle subselects
* handle aggregates on views
  * handle insert into select from view works
System indexes are now multikey(Bruce)
Oidint2, oidint4, and oidname types are removed(Bruce)
Use system cache for more system table lookups(Bruce)
New backend programming language PL/pgSQL in backend/pl(Jan)
New SERIAL data type, auto-creates sequence/index(Thomas)
Enable assert checking without a recompile(Massimo)
User lock enhancements(Massimo)
New setval() command to set sequence value(Massimo)
Auto-remove unix socket file on start-up if no postmaster running(Massimo)
Conditional trace package(Massimo)
New UNLISTEN command(Massimo)
psql and libpq now compile under Windows using win32.mak(Magnus)
Lo_read no longer stores trailing NULL(Bruce)
Identifiers are now truncated to 31 characters internally(Bruce)
Createuser options now available on the command line
Code for 64-bit integer supported added, configure tested, int8 type(Thomas)
Prevent file descriptor leaf from failed COPY(Bruce)
New pg_upgrade command(Bruce)
Updated /contrib directories(Massimo)
New CREATE TABLE DEFAULT VALUES statement available(Thomas)
New INSERT INTO TABLE DEFAULT VALUES statement available(Thomas)
New DECLARE and FETCH feature(Thomas)
libpq's internal structures now not exported(Tom)
Allow up to 8 key indexes(Bruce)
Remove ARCHIVE key word, that is no longer used(Thomas)
pq_dump -n flag to suppress quotes around identifiers
disable system columns for views(Jan)
new INET and CIDR types for network addresses(TomH, Paul)
no more double quotes in psql output
pq_dump now dumps views(Terry)
new SET QUERY_LIMIT(Tatsuo,Jan)

Source Tree Changes
-------------------
/contrib cleanup(Jun)
Inline some small functions called for every row(Bruce)
Alpha/linux fixes
Release Notes

HP-UX cleanups (Tom)
Multibyte regression tests (Soonmyung.)
Remove --disabled options from configure
Define PGDOC to use POSTGRESDIR by default
Make regression optional
Remove extra braces code to pgindent (Bruce)
Add bsd shared library support (Bruce)
New --without-CXX support configure option (Brook)
New FAQ_CVS
Update backend flowchart in tools/backend (Bruce)
Change atttypmod from int16 to int32 (Bruce, Tom)
Getrusage() fix for platforms that do not have it (Tom)
Add PQconnectdb, PGUSER, PGPASSWORD to libpq man page
NS32K platform fixes (Phil Nelson, John Buller)
SCO 7/UnixWare 2.x fixes (Billy, others)
Sparc/Solaris 2.5 fixes (Ryan)
Pgbuiltin.3 is obsolete, move to doc files (Thomas)
Even more documentation (Thomas)
Nextstep support (Jacek)
Aix support (David)
pginterface manual page (Bruce)
shared libraries all have version numbers
merged all OS-specific shared library defines into one file
smarter TCL/TK configuration checking (Billy)
smarter perl configuration (Brook)
configure uses supplied install-sh if no install script found (Tom)
new Makefile.shlib for shared library configuration (Tom)

E.358. Release 6.3.2

Release date: 1998-04-07

This is a bug-fix release for 6.3.x. Refer to the release notes for version 6.3 for a more complete
summary of new features.

Summary:

- Repairs automatic configuration support for some platforms, including Linux, from breakage inad-
  vertently introduced in version 6.3.1.

- Correctly handles function calls on the left side of BETWEEN and LIKE clauses.

A dump/restore is NOT required for those running 6.3 or 6.3.1. A make distclean, make, and
make install is all that is required. This last step should be performed while the postmaster is not
running. You should re-link any custom applications that use PostgreSQL libraries.

For upgrades from pre-6.3 installations, refer to the installation and migration instructions for version
6.3.

E.358.1. Changes

Configure detection improvements for tcl/tk (Brook Milligan, Alvin)
Manual page improvements (Bruce)
BETWEEN and LIKE fix (Thomas)
fix for psql \connect used by pg_dump (Oliver Elphick)
New odbc driver
pgaccess, version 0.86
qsort removed, now uses libc version, cleanups (Jeroen)
Release Notes

fix for buffer over-runs detected (Maurice Gittens)
fix for buffer overrun in libpqctl (Randy Kunkee)
fix for UNION with DISTINCT or ORDER BY (Bruce)
gettimeofday configure check (Doug Winterburn)
Fix "indexes not used" bug (Vadim)
docs additions (Thomas)
Fix for backend memory leak (Bruce)
libreadline cleanup (Erwan MAS)
Remove DISTDIR (Bruce)
Makefile dependency cleanup (Jeroen van Vianen)
ASSERT fixes (Bruce)

E.359. Release 6.3.1

Release date: 1998-03-23

Summary:

• Additional support for multibyte character sets.
• Repair byte ordering for mixed-endian clients and servers.
• Minor updates to allowed SQL syntax.
• Improvements to the configuration autodetection for installation.

A dump/restore is NOT required for those running 6.3. A make distclean, make, and make install is all that is required. This last step should be performed while the postmaster is not running. You should re-link any custom applications that use PostgreSQL libraries.

For upgrades from pre-6.3 installations, refer to the installation and migration instructions for version 6.3.

E.359.1. Changes

ecpg cleanup/fixes, now version 1.1 (Michael Meskes)
pg_user cleanup (Bruce)
large object fix for pg_dump and tclsh (alvin)
LIKE fix for multiple adjacent underscores
fix for redefining builtin functions (Thomas)
ultrix4 cleanup
upgrade to pg_access 0.83
updated CLUSTER manual page
multibyte character set support, see doc/README.mb (Tatsuo)
configure --with-pgport fix
pg_ident fix
big-endian fix for backend communications (Kataoka)
SUBSTR() and substring() fix (Jan)
several jdbc fixes (Peter)
libpqctl improvements, see libpqctl/README (Randy Kunkee)
Fix for "Datasize = 0" error (Vadim)
Prevent \do from wrapping (Bruce)
Remove duplicate Russian character set entries
Sunos4 cleanup
Allow optional TABLE key word in LOCK and SELECT INTO (Thomas)
CREATE SEQUENCE options to allow a negative integer (Thomas)
Add "PASSWORD" as an allowed column identifier (Thomas)
Add checks for UNION target fields (Bruce)
E.360. Release 6.3

Release date: 1998-03-01

There are many new features and improvements in this release. Here is a brief, incomplete summary:

• Many new SQL features, including full SQL92 subselect capability (everything is here but target-list subselects).

• Support for client-side environment variables to specify time zone and date style.

• Socket interface for client/server connection. This is the default now so you might need to start postmaster with the -i flag.

• Better password authorization mechanisms. Default table privileges have changed.

• Old-style time travel has been removed. Performance has been improved.

Note

Bruce Momjian wrote the following notes to introduce the new release.

There are some general 6.3 issues that I want to mention. These are only the big items that cannot be described in one sentence. A review of the detailed changes list is still needed.

First, we now have subselects. Now that we have them, I would like to mention that without subselects, SQL is a very limited language. Subselects are a major feature, and you should review your code for places where subselects provide a better solution for your queries. I think you will find that there are more uses for subselects than you might think. Vadim has put us on the big SQL map with subselects, and fully functional ones too. The only thing you cannot do with subselects is to use them in the target list.

Second, 6.3 uses Unix domain sockets rather than TCP/IP by default. To enable connections from other machines, you have to use the new postmaster -i option, and of course edit pg_hba.conf. Also, for this reason, the format of pg_hba.conf has changed.

Third, char() fields will now allow faster access than varchar() or text. Specifically, the text and varchar() have a penalty for access to any columns after the first column of this type. char() used to also have this access penalty, but it no longer does. This might suggest that you redesign some of your tables, especially if you have short character columns that you have defined as varchar() or text. This and other changes make 6.3 even faster than earlier releases.

We now have passwords definable independent of any Unix file. There are new SQL USER commands. See the Administrator's Guide for more information. There is a new table, pg_shadow, which is used to store user information and user passwords, and it by default only SELECT-able by the postgres super-user. pg_user is now a view of pg_shadow, and is SELECT-able by PUBLIC. You should keep using pg_user in your application without changes.

User-created tables now no longer have SELECT privilege to PUBLIC by default. This was done because the ANSI standard requires it. You can of course GRANT any privileges you want after the table is created. System tables continue to be SELECT-able by PUBLIC.

We also have real deadlock detection code. No more sixty-second timeouts. And the new locking code implements a FIFO better, so there should be less resource starvation during heavy use.
Many complaints have been made about inadequate documentation in previous releases. Thomas has put much effort into many new manuals for this release. Check out the doc/ directory.

For performance reasons, time travel is gone, but can be implemented using triggers (see `pgsql/contrib/spi/README`). Please check out the new \d command for types, operators, etc. Also, views have their own privileges now, not based on the underlying tables, so privileges on them have to be set separately. Check `/pgsql/interfaces` for some new ways to talk to PostgreSQL.

This is the first release that really required an explanation for existing users. In many ways, this was necessary because the new release removes many limitations, and the work-arounds people were using are no longer needed.

### E.360.1. Migration to Version 6.3

A dump/restore using `pg_dump` or `pg_dumpall` is required for those wishing to migrate data from any previous release of PostgreSQL.

### E.360.2. Changes

#### Bug Fixes

-----

- Fix binary cursors broken by MOVE implementation (Vadim)
- Fix for tcl library crash (Jan)
- Fix for array handling, from Gerhard Hintermayer
- Fix acl error, and remove duplicate pgtrace (Bruce)
- Fix psql `\e` for empty file (Bruce)
- Fix for textcat on varchar() fields (Bruce)
- Fix for DBT Sendproc (Zeugswetter Andres)
- Fix vacuum analyze syntax problem (Bruce)
- Fix for international identifiers (Tatsuo)
- Fix aggregates on inherited tables (Bruce)
- Fix `substr()` for out-of-bounds data
- Fix for select `1=1` or `2=2`, select `1=1` and `2=2`, and select `sum(2+2)` (Bruce)
- Fix notty output to show status result. `-q` option still turns it off (Bruce)
- Fix for `count(*)`, aggs with views and multiple tables and `sum(3)` (Bruce)
- Fix cluster (Bruce)
- Fix for PQtrace start/stop several times (Bruce)
- Fix a variety of locking problems like newer lock waiters getting lock before older waiters, and having readlock people not share locks if a writer is waiting for a lock, and waiting writers not getting priority over waiting readers (Bruce)
- Fix crashes in psql when executing queries from external files (James)
- Fix problem with multiple order by columns, with the first one having NULL values (Jeroen)
- Use correct hash table support functions for float8 and int4 (Thomas)
- Re-enable JOIN= option in CREATE OPERATOR statement (Thomas)
- Change precedence for boolean operators to match expected behavior (Thomas)
- Generate elog(ERROR) on over-large integer (Bruce)
Release Notes

Allow multiple-argument functions in constraint clauses (Thomas)
Check boolean input literals for 'true', 'false', 'yes', 'no', '1', '0'
and throw elog(ERROR) if unrecognized (Thomas)

Major large objects fix
Fix for GROUP BY showing duplicates (Vadim)
Fix for index scans in MergeJoin (Vadim)

Enhancements

Subselects with EXISTS, IN, ALL, ANY key words (Vadim, Bruce, Thomas)
New User Manual (Thomas, others)
Speedup by inlining some frequently-called functions
Real deadlock detection, no more timeouts (Bruce)
Add SQL92 "constants" CURRENT_DATE, CURRENT_TIME,
   CURRENT_TIMESTAMP,
   CURRENT_USER (Thomas)
Modify constraint syntax to be SQL92-compliant (Thomas)
Implement SQL92 PRIMARY KEY and UNIQUE clauses using indexes (Thomas)
Recognize SQL92 syntax for FOREIGN KEY. Throw elog notice (Thomas)
Allow NOT NULL UNIQUE constraint clause (each allowed separately before) (Thomas)
Allow PostgreSQL-style casting ("::") of non-constants (Thomas)
Add support for SQL3 TRUE and FALSE boolean constants (Thomas)
Support SQL92 syntax for IS TRUE/IS FALSE/IS NOT TRUE/IS NOT FALSE (Thomas)
Allow shorter strings for boolean literals (e.g. "t", "tr", "tru") (Thomas)
Allow SQL92 delimited identifiers (Thomas)
Implement SQL92 binary and hexadecimal string decoding (b'10' and x'1F') (Thomas)
Support SQL92 syntax for type coercion of literal strings
   (e.g. "DATETIME 'now'") (Thomas)
Add conversions for int2, int4, and OID types to and from text (Thomas)
Use shared lock when building indexes (Vadim)
Free memory allocated for a user query inside transaction block
   after this query is done, was turned off in <= 6.2.1 (Vadim)
New SQL statement CREATE PROCEDURAL LANGUAGE (Jan)
New PostgreSQL procedural language (PL) backend interface (Jan)
Rename pg_dump -H option to -h (Bruce)
Add Java support for passwords, European dates (Peter)
Use indexes for LIKE and ~, !~ operations (Bruce)
Add hash functions for datetime and timespan (Thomas)
Time Travel removed (Vadim, Bruce)
Add paging for \d and \z, and fix \i (Bruce)
Add Unix domain socket support to backend and to frontend
   library (Goran)
Implement CREATE DATABASE/WITH LOCATION and initlocation
   utility (Thomas)
Allow more SQL92 and/or PostgreSQL reserved words as column
   identifiers (Thomas)
Augment support for SQL92 SET TIME ZONE... (Thomas)
SET/SHOW/RESET TIME ZONE uses TZ backend environment
   variable (Thomas)
Implement SET keyword = DEFAULT and SET TIME ZONE DEFAULT (Thomas)
Enable SET TIME ZONE using TZ environment variable (Thomas)
Add PGDATESTYLE environment variable to frontend and backend initialization (Thomas)
Add PGTZ, PG_COST_HEAP, PG_COST_INDEX, PGRPLANS, PG_GEQO
  frontend library initialization environment variables (Thomas)
Regression tests time zone automatically set with "setenv PGTZ PST8PDT" (Thomas)
Add pg_description table for info on tables, columns, operators, types, and aggregates (Bruce)
Increase 16 char limit on system table/index names to 32 characters (Bruce)
Rename system indexes (Bruce)
Add 'GERMAN' option to SET DATESTYLE (Thomas)
Define an "ISO-style" timespan output format with "hh:mm:ss" fields (Thomas)
Allow fractional values for delta times (e.g. '2.5 days') (Thomas)
Validate numeric input more carefully for delta times (Thomas)
Implement day of year as possible input to date_part() (Thomas)
Define timespan_finite() and text_timespan() functions (Thomas)
Remove archive stuff (Bruce)
Allow for a pg_password authentication database that is separate from the system password file (Todd)
Dump ACLs, GRANT, REVOKE privileges (Matt)
Define text, varchar, and bpchar string length functions (Thomas)
Fix Query handling for inheritance, and cost computations (Bruce)
Implement CREATE TABLE/AS SELECT (alternative to SELECT/INTO) (Thomas)
Allow NOT, IS NULL, IS NOT NULL in constraints (Thomas)
Implement UNIONs for SELECT (Bruce)
Add UNION, GROUP, DISTINCT to INSERT (Bruce)
varchar() stores only necessary bytes on disk (Bruce)
Fix for BLOBs (Peter)
Mega-Patch for JDBC...see README_6.3 for list of changes (Peter)
Remove unused "option" from PQconnectdb() New LOCK command and lock manual page describing deadlocks (Bruce)
Add new psql \da, \dd, \df, \do, \ds, and \dT commands (Bruce)
Enhance psql \z to show sequences (Bruce)
Show NOT NULL and DEFAULT in psql \d table (Bruce)
New psql .psqlrc file start-up (Andrew)
Modify sample start-up script in contrib/linux to show syslog (Thomas)
New types for IP and MAC addresses in contrib/ip_and_mac (TomH)
Unix system time conversions with date/time types in contrib/unixdate (Thomas)
Update of contrib stuff (Massimo)
Add Unix socket support to DBD::Pg (Goran)
New python interface (PyGreSQL 2.0) (D'Arcy)
New frontend/backend protocol has a version number, network byte order (Phil)
Security features in pg_hba.conf enhanced and documented, many cleanups (Phil)
CHAR() now faster access than VARCHAR() or TEXT
ecpg embedded SQL preprocessor
Reduce system column overhead (Vadmin)
Remove pg_time table (Vadim)
Add pg_type attribute to identify types that need length (bpchar, varchar)
Add report of offending line when COPY command fails
Allow VIEW privileges to be set separately from the underlying tables.
   For security, use GRANT/REVOKE on views as appropriate (Jan)
Tables now have no default GRANT SELECT TO PUBLIC. You must explicitly grant such privileges.
Clean up tutorial examples (Darren)

Source Tree Changes
-------------------
Add new html development tools, and flow chart in /tools/backend
Fix for SCO compiles
Stratus computer port Robert Gillies
Added support for shlib for BSD44-derived & i386_solaris
Make configure more automated (Brook)
Add script to check regression test results
Break parser functions into smaller files, group together (Bruce)
Rename heap_create to heap_create_and_catalog, rename heap_creatr to heap_create() (Bruce)
Sparc/Linux patch for locking (TomS)
Remove PORTNAME and reorganize port-specific stuff (Marc)
Add optimizer README file (Bruce)
Remove some recursion in optimizer and clean up some code there (Bruce)
Fix for NetBSD locking (Henry)
Fix for libptcl make (Tatsuo)
AIX patch (Darren)
Change IS TRUE, IS FALSE, ... to expressions using "=" rather than function calls to istrue() or isfalse() to allow optimization (Thomas)
Various fixes NetBSD/Sparc related (TomH)
Alpha linux locking (Travis, Ryan)
Change elog(WARN) to elog(ERROR) (Bruce)
FAQ for FreeBSD (Marc)
Bring in the PostODBC source tree as part of our standard distribution (Marc)
A minor patch for HP/UX 10 vs 9 (Stan)
New pg_attribute.atttypmod for type-specific info like varchar length (Bruce)
UnixWare patches (Billy)
New i386 'lock' for spinlock asm (Billy)
Support for multiplexed backends is removed
Start an OpenBSD port
Start an AUX port
Start a Cygnus port
Add string functions to regression suite (Thomas)
Expand a few function names formerly truncated to 16 characters (Thomas)
Remove un-needed malloc() calls and replace with palloc() (Bruce)

E.361. Release 6.2.1

Release date:  1997-10-17

6.2.1 is a bug-fix and usability release on 6.2.
Summary:

• Allow strings to span lines, per SQL92.
• Include example trigger function for inserting user names on table updates.

This is a minor bug-fix release on 6.2. For upgrades from pre-6.2 systems, a full dump/reload is required. Refer to the 6.2 release notes for instructions.

E.361.1. Migration from version 6.2 to version 6.2.1

This is a minor bug-fix release. A dump/reload is not required from version 6.2, but is required from any release prior to 6.2.

In upgrading from version 6.2, if you choose to dump/reload you will find that avg(money) is now calculated correctly. All other bug fixes take effect upon updating the executables.

Another way to avoid dump/reload is to use the following SQL command from psql to update the existing system table:

```
update pg_aggregate set aggfinalfn = 'cash_div_flt8'
where aggnname = 'avg' and aggbasetype = 790;
```

This will need to be done to every existing database, including template1.

E.361.2. Changes

Allow TIME and TYPE column names (Thomas)
Allow larger range of true/false as boolean values (Thomas)
Support output of "now" and "current" (Thomas)
Handle DEFAULT with INSERT of NULL properly (Vadim)
Fix for relation reference counts problem in buffer manager (Vadim)
Allow strings to span lines, like ANSI (Thomas)
Fix for backward cursor with ORDER BY (Vadim)
Fix avg(cash) computation (Thomas)
Fix for specifying a column twice in ORDER/GROUP BY (Vadim)
Documented new libpq function to return affected rows, PQcmdTuples (Bruce)
Trigger function for inserting user names for INSERT/UPDATE (Brook Milligan)

E.362. Release 6.2

Release date: 1997-10-02

A dump/restore is required for those wishing to migrate data from previous releases of PostgreSQL.

E.362.1. Migration from version 6.1 to version 6.2

This migration requires a complete dump of the 6.1 database and a restore of the database in 6.2.

Note that the pg_dump and pg_dumpall utility from 6.2 should be used to dump the 6.1 database.

E.362.2. Migration from version 1.* to version 6.2

Those migrating from earlier 1.* releases should first upgrade to 1.09 because the COPY output format was improved from the 1.02 release.
E.362.3. Changes

Bug Fixes
--------
Fix problems with pg_dump for inheritance, sequences, archive tables (Bruce)
Fix compile errors on overflow due to shifts, unsigned, and bad prototypes
from Solaris (Diab Jerius)
Fix bugs in geometric line arithmetic (bad intersection calculations) (Thomas)
Check for geometric intersections at endpoints to avoid rounding ugliness (Thomas)
Catch non-functional delete attempts (Vadim)
Change time function names to be more consistent (Michael Reifenberg)
Check for zero divides (Michael Reifenberg)
Fix very old bug which made rows changed/inserted by a command visible to the command itself (so we had multiple update of updated rows, etc.) (Vadim)
Fix for SELECT null, 'fail' FROM pg_am (Patrick)
SELECT NULL as EMPTY_FIELD now allowed (Patrick)
Remove un-needed signal stuff from contrib/pginterface
Fix OR (where x != 1 or x isnull didn't return rows with x NULL) (Vadim)
Fix time_cmp function (Vadim)
Fix handling of functions with non-attribute first argument in WHERE clauses (Vadim)
Fix GROUP BY when order of entries is different from order in target list (Vadim)
Fix pg_dump for aggregates without sfunc1 (Vadim)

Enhancements
----------
Default genetic optimizer GEQO parameter is now 8 (Bruce)
Allow use parameters in target list having aggregates in functions (Vadim)
Added JDBC driver as an interface (Adrian & Peter)
pg_password utility
Return number of rows inserted/affected by INSERT/UPDATE/DELETE etc. (Vadim)
Triggers implemented with CREATE TRIGGER (SQL3) (Vadim)
SPI (Server Programming Interface) allows execution of queries inside C-functions (Vadim)
NOT NULL implemented (SQL92) (Robson Paniago de Miranda)
Include reserved words for string handling, outer joins, and unions (Thomas)
Implement extended comments ("/* ... */") using exclusive states (Thomas)
Add "//" single-line comments (Bruce)
Remove some restrictions on characters in operator names (Thomas)
DEFAULT and CONSTRAINT for tables implemented (SQL92) (Vadim & Thomas)
Add text concatenation operator and function (SQL92) (Thomas)
Support WITH TIME ZONE syntax (SQL92) (Thomas)
Support INTERVAL unit TO unit syntax (SQL92) (Thomas)
Define types DOUBLE PRECISION, INTERVAL, CHARACTER, and CHARACTER VARYING (SQL92) (Thomas)
Define type FLOAT(p) and rudimentary DECIMAL(p,s), NUMERIC(p,s) (SQL92) (Thomas)
Define EXTRACT(), POSITION(), SUBSTRING(), and TRIM() (SQL92) (Thomas)
Define CURRENT_DATE, CURRENT_TIME, CURRENT_TIMESTAMP (SQL92) (Thomas)
Add syntax and warnings for UNION, HAVING, INNER and OUTER JOIN (SQL92) (Thomas)
Add more reserved words, mostly for SQL92 compliance (Thomas)
Allow hh:mm:ss time entry for timespan/reltime types (Thomas)
Add center() routines for lseg, path, polygon (Thomas)
Add distance() routines for circle-polygon, polygon-polygon (Thomas)
Check explicitly for points and polygons contained within polygons using an axis-crossing algorithm (Thomas)
Add routine to convert circle-box (Thomas)
Merge conflicting operators for different geometric data types (Thomas)
Replace distance operator "<===" with "<->" (Thomas)
Replace "above" operator "!^" with "^^" and "below" operator "!|" with "<^" (Thomas)
Add routines for text trimming on both ends, substring, and string position (Thomas)
Added conversion routines circle(box) and poly(circle) (Thomas)
Allow internal sorts to be stored in memory rather than in files (Bruce & Vadim)
Allow functions and operators on internally-identical types to succeed (Bruce)
Speed up backend start-up after profiling analysis (Bruce)
Inline frequently called functions for performance (Bruce)
Reduce open() calls (Bruce)
psql: Add PAGER for \h and \?,\C fix
Fix for psql pager when no tty (Bruce)
New entab utility (Bruce)
General trigger functions for referential integrity (Vadim)
General trigger functions for time travel (Vadim)
General trigger functions for AUTOINCREMENT/IDENTITY feature (Vadim)
MOVE implementation (Vadim)

Source Tree Changes
-------------------
HP-UX 10 patches (Vladimir Turin)
Added SCO support, (Daniel Harris)
MkLinux patches (Tatsuo Ishii)
Change geometric box terminology from "length" to "width" (Thomas)
Deprecate temporary unstored slope fields in geometric code (Thomas)
Remove restart instructions from INSTALL (Bruce)
Look in /usr/ucb first for install (Bruce)
Fix c++ copy example code (Thomas)
Add -o to psql manual page (Bruce)
Prevent rename unallocated string length from being copied into database (Bruce)
Cleanup for NAMEDATALEN use (Bruce)
Fix pg_proc names over 15 chars in output (Bruce)
Add strNcpy() function (Bruce)
remove some (void) casts that are unnecessary (Bruce)
new interfaces directory (Marc)
Replace fopen() calls with calls to fd.c functions (Bruce)
Make functions static where possible (Bruce)
enclose unused functions in #ifdef NOT_USED (Bruce)
Remove call to difftime() in timestamp support to fix SunOS (Bruce & Thomas)
Changes for Digital Unix
Portability fix for pg_dumpall (Bruce)
Rename pg_attribute.attnvals to attdispersion (Bruce)
"intro/unix" manual page now "pgintro" (Bruce)
"built-in" manual page now "pgbuiltin" (Bruce)
"drop" manual page now "drop_table" (Bruce)
Add "create_trigger", "drop_trigger" manual pages (Thomas)
Add constraints regression test (Vadim & Thomas)
Add comments syntax regression test (Thomas)
Add PGINDENT and support program (Bruce)
Massive commit to run PGINDENT on all *.c and *.h files (Bruce)
Files moved to /src/tools directory (Bruce)
SPI and Trigger programming guides (Vadim & D'Arcy)

E.363. Release 6.1.1

Release date: 1997-07-22

E.363.1. Migration from version 6.1 to version 6.1.1

This is a minor bug-fix release. A dump/reload is not required from version 6.1, but is required from any release prior to 6.1. Refer to the release notes for 6.1 for more details.

E.363.2. Changes

fix for SET with options (Thomas)
allow pg_dump/pg_dumpall to preserve ownership of all tables/objects (Bruce)
new psql \connect option allows changing usernames without changing databases
fix for initdb --debug option (Yoshihiko Ichikawa)
lextest cleanup (Bruce)
hash fixes (Vadim)
fix date/time month boundary arithmetic (Thomas)
fix timezone daylight handling for some ports (Thomas, Bruce, Tatsuo)
timestamp overhauled to use standard functions (Thomas)
other code cleanup in date/time routines (Thomas)
psql's \d now case-insensitive (Bruce)
psql's backslash commands can now have trailing semicolon (Bruce)
fix memory leak in psql when using \g (Bruce)
major fix for endian handling of communication to server (Thomas, Tatsuo)
Fix for Solaris assembler and include files (Yoshihiko Ichikawa)
allow underscores in usernames (Bruce)
pg_dumpall now returns proper status, portability fix (Bruce)
E.364. Release 6.1

**Release date:** 1997-06-08

The regression tests have been adapted and extensively modified for the 6.1 release of PostgreSQL.

Three new data types (datetime, timespan, and circle) have been added to the native set of PostgreSQL types. Points, boxes, paths, and polygons have had their output formats made consistent across the data types. The polygon output in misc.out has only been spot-checked for correctness relative to the original regression output.

PostgreSQL 6.1 introduces a new, alternate optimizer which uses genetic algorithms. These algorithms introduce a random behavior in the ordering of query results when the query contains multiple qualifiers or multiple tables (giving the optimizer a choice on order of evaluation). Several regression tests have been modified to explicitly order the results, and hence are insensitive to optimizer choices. A few regression tests are for data types which are inherently unordered (e.g. points and time intervals) and tests involving those types are explicitly bracketed with `set geqo to 'off'` and `reset geqo`.

The interpretation of array specifiers (the curly braces around atomic values) appears to have changed sometime after the original regression tests were generated. The current `.expected/*.out` files reflect this new interpretation, which might not be correct!

The float8 regression test fails on at least some platforms. This is due to differences in implementations of `pow()` and `exp()` and the signaling mechanisms used for overflow and underflow conditions.

The “random” results in the random test should cause the “random” test to be “failed”, since the regression tests are evaluated using a simple diff. However, “random” does not seem to produce random results on my test machine (Linux/gcc/i686).

### E.364.1. Migration to Version 6.1

This migration requires a complete dump of the 6.0 database and a restore of the database in 6.1.

Those migrating from earlier 1.* releases should first upgrade to 1.09 because the COPY output format was improved from the 1.02 release.

### E.364.2. Changes

**Bug Fixes**

- packet length checking in library routines
- lock manager priority patch
- check for under/over flow of float8 (Bruce)
- multitable join fix (Vadim)
- SIGPIPE crash fix (Darren)
- large object fixes (Sven)
- allow btree indexes to handle NULLs (Vadim)
- timezone fixes (D'Arcy)
- select SUM(x) can return NULL on no rows (Thomas)
- internal optimizer, executor bug fixes (Vadim)
- fix problem where inner loop in < or <= has no rows (Vadim)
- prevent re-commuting join index clauses (Vadim)
- fix join clauses for multiple tables (Vadim)
- fix hash, hashjoin for arrays (Vadim)
- fix btree for abstime type (Vadim)
- large object fixes (Raymond)
- fix buffer leak in hash indexes (Vadim)
- fix rtree for use in inner scan (Vadim)
Release Notes

fix gist for use in inner scan, cleanups (Vadim, Andrea)
avoid unnecessary local buffers allocation (Vadim, Massimo)
fix local buffers leak in transaction aborts (Vadim)
fix file manager memory leaks, cleanups (Vadim, Massimo)
fix storage manager memory leaks (Vadim)
fix btree duplicates handling (Vadim)
fix deleted rows reincarnation caused by vacuum (Vadim)
fix SELECT varchar()/char() INTO TABLE made zero-length fields (Bruce)
many psql, pg_dump, and libpq memory leaks fixed using Purify (Igor)

Enhancements
-------------
attribute optimization statistics (Bruce)
much faster new btree bulk load code (Paul)
BTREE UNIQUE added to bulk load code (Vadim)
new lock debug code (Massimo)
massive changes to libpq++ (Leo)
new GEQO optimizer speeds table multitable optimization (Martin)
new WARN message for non-unique insert into unique key (Marc)
update x=-3, no spaces, now valid (Bruce)
remove case-sensitive identifier handling (Bruce, Thomas, Dan)
debug backend now pretty-prints tree (Darren)
new Oracle character functions (Edmund)
new plaintext password functions (Dan)
no such class or insufficient privilege changed to distinct messages (Dan)
new ANSI timestamp function (Dan)
new ANSI Time and Date types (Thomas)
move large chunks of data in backend (Martin)
multicolumn btree indexes (Vadim)
new SET var TO value command (Martin)
update transaction status on reads (Dan)
new locale settings for character types (Oleg)
new SEQUENCE serial number generator (Vadim)
GROUP BY function now possible (Vadim)
re-organize regression test (Thomas, Marc)
new optimizer operation weights (Vadim)
new psql \z grant/permit option (Marc)
new MONEY data type (D'Arcy, Thomas)
tcp socket communication speed improved (Vadim)
new VACUUM option for attribute statistics, and for certain columns (Vadim)
many geometric type improvements (Thomas, Keith)
additional regression tests (Thomas)
new datestyle variable (Thomas, Vadim, Martin)
more comparison operators for sorting types (Thomas)
new conversion functions (Thomas)
new more compact btree format (Vadim)
allow pg_dumpall to preserve database ownership (Bruce)
new SET GEQO=# and R_PLANS variable (Vadim)
old (!GEQO) optimizer can use right-sided plans (Vadim)
typechecking improvement in SQL parser (Bruce)
new SET, SHOW, RESET commands (Thomas, Vadim)
new \connect database USER option
new destroydb -i option (Igor)
new \dt and \di psql commands (Darren)
SELECT "\n" now escapes newline (A. Duursma)
new geometry conversion functions from old format (Thomas)

Source tree changes
-------------------
new configuration script (Marc)
readline configuration option added (Marc)
OS-specific configuration options removed (Marc)
new OS-specific template files (Marc)
no more need to edit Makefile.global (Marc)
re-arrange include files (Marc)
exterstep patches (Gregor Hoffleit)
removed Windows-specific code (Bruce)
removed postmaster -e option, now only postgres -e option (Bruce)
merge duplicate library code in front/backends (Martin)
now works with eBones, international Kerberos (Jun)
more shared library support
C++ include file cleanup (Bruce)
warn about buggy flex (Bruce)
DG/UX, Ultrix, IRIX, AIX portability fixes

E.365. Release 6.0

Release date: 1997-01-29

A dump/restore is required for those wishing to migrate data from previous releases of PostgreSQL.

E.365.1. Migration from version 1.09 to version 6.0

This migration requires a complete dump of the 1.09 database and a restore of the database in 6.0.

E.365.2. Migration from pre-1.09 to version 6.0

Those migrating from earlier 1.* releases should first upgrade to 1.09 because the COPY output format was improved from the 1.02 release.

E.365.3. Changes

Bug Fixes
--------
ALTER TABLE bug - running postgres process needs to re-read table definition
Allow vacuum to be run on one table or entire database (Bruce)
Array fixes
Fix array over-runs of memory writes (Kurt)
Fix elusive btree range/non-range bug (Dan)
Fix for hash indexes on some types like time and date
Fix for pg_log size explosion
Fix permissions on lo_export () (Bruce)
Fix uninitialized reads of memory (Kurt)
Fixed ALTER TABLE ... char(3) bug (Bruce)
Fixed a few small memory leaks
Fixed EXPLAIN handling of options and changed full_path option name
Fixed output of group acl privileges
Memory leaks (hunt and destroy with tools like Purify (Kurt)
Minor improvements to rules system
NOTIFY fixes
New asserts for run-checking
Overhauled parser/analyze code to properly report errors and increase speed
Pg_dump -d now handles NULL's properly (Bruce)
Prevent SELECT NULL from crashing server (Bruce)
Properly report errors when INSERT ... SELECT columns did not match
Properly report errors when insert column names were not correct
psql \g filename now works (Bruce)
psql fixed problem with multiple statements on one line with multiple outputs
Removed duplicate system OIDs
SELECT * INTO TABLE . GROUP/ORDER BY gives unlink error if table exists (Bruce)
Several fixes for queries that crashed the backend
Starting quote in insert string errors (Bruce)
Submitting an empty query now returns empty status, not just " " query (Bruce)

Enhancements
------------
Add EXPLAIN manual page (Bruce)
Add UNIQUE index capability (Dan)
Add hostname/user level access control rather than just hostname and user
Add synonym of != for <>(Bruce)
Allow "select oid,* from table"
Allow BY,ORDER BY to specify columns by number, or by non-alias table.column (Bruce)
Allow COPY from the frontend (Bryan)
Allow GROUP BY to use alias column name (Bruce)
Allow actual compression, not just reuse on the same page (Vadim)
Allow installation-configuration option to auto-add all local users (Bryan)
Allow libpq to distinguish between text value '' and null (Bruce)
Allow non-postgres users with createdb privs to destroydb's
Allow restriction on who can create C functions (Bryan)
Allow restriction on who can do backend COPY (Bryan)
Can shrink tables, pg_time and pg_log (Vadim & Erich)
Change debug level 2 to print queries only, changed debug heading layout (Bruce)
Change default decimal constant representation from float4 to float8 (Bruce)
European date format now set when postmaster is started
Execute lowercase function names if not found with exact case
Fixes for aggregate/GROUP processing, allow 'select sum(func(x),sum(x+y) from z'
Gist now included in the distribution (Marc)
Ident authentication of local users (Bryan)
Implement BETWEEN qualifier (Bruce)
Implement IN qualifier (Bruce)
libpq has PQgetisnull() (Bruce)
libpq++ improvements
New options to initdb (Bryan)
Pg_dump allow dump of OIDs (Bruce)
Pg_dump create indexes after tables are loaded for speed (Bruce)
Pg_dumpall dumps all databases, and the user table
Pginterface additions for NULL values (Bruce)
Prevent postmaster from being run as root
psql \h and \? is now readable (Bruce)
psql allow backslashed, semicolons anywhere on the line (Bruce)
psql changed command prompt for lines in query or in quotes (Bruce)
psql char(3) now displays as (bp)char in \d output (Bruce)
psql return code now more accurate (Bryan?)
psql updated help syntax (Bruce)
Re-visit and fix vacuum (Vadim)
Reduce size of regression diffs, remove timezone name
difference (Bruce)
Remove compile-time parameters to enable binary
distributions (Bryan)
Reverse meaning of HBA masks (Bryan)
Secure Authentication of local users (Bryan)
Speed up vacuum (Vadim)
Vacuum now had VERBOSE option (Bruce)

Source tree changes
-------------------
All functions now have prototypes that are compared against the calls
Allow asserts to be disabled easily from Makefile.global (Bruce)
Change oid constants used in code to #define names
Decoupled sparc and solaris defines (Kurt)
Gcc -Wall compiles cleanly with warnings only from unfixable constructs
Major include file reorganization/reduction (Marc)
Make now stops on compile failure (Bryan)
Makefile restructuring (Bryan, Marc)
Merge bsdi_2.1 to bsdi (Bruce)
Monitor program removed
Name change from Postgres95 to PostgreSQL
New config.h file (Marc, Bryan)
PG_VERSION now set to 6.0 and used by postmaster
Portability additions, including Ultrix, DG/UX, AIX, and Solaris
Reduced the number of #define's, centralized #define's
Remove duplicate OIDS in system tables (Dan)
Remove duplicate system catalog info or report mismatches (Dan)
Removed many os-specific #define's
Restructured object file generation/location (Bryan, Marc)
Restructured port-specific file locations (Bryan, Marc)
Unused/uninitialized variables corrected

E.366. Release 1.09

Release date: 1996-11-04

Sorry, we didn't keep track of changes from 1.02 to 1.09. Some of the changes listed in 6.0 were actually included in the 1.02.1 to 1.09 releases.

E.367. Release 1.02

Release date: 1996-08-01
E.367.1. Migration from version 1.02 to version 1.02.1

Here is a new migration file for 1.02.1. It includes the ‘copy’ change and a script to convert old ASCII files.

Note

The following notes are for the benefit of users who want to migrate databases from Postgres95 1.01 and 1.02 to Postgres95 1.02.1.

If you are starting afresh with Postgres95 1.02.1 and do not need to migrate old databases, you do not need to read any further.

In order to upgrade older Postgres95 version 1.01 or 1.02 databases to version 1.02.1, the following steps are required:

1. Start up a new 1.02.1 postmaster
2. Add the new built-in functions and operators of 1.02.1 to 1.01 or 1.02 databases. This is done by running the new 1.02.1 server against your own 1.01 or 1.02 database and applying the queries attached at the end of the file. This can be done easily through `psql`. If your 1.01 or 1.02 database is named `testdb` and you have cut the commands from the end of this file and saved them in `addfunc.sql`:

```
% psql testdb -f addfunc.sql
```

Those upgrading 1.02 databases will get a warning when executing the last two statements in the file because they are already present in 1.02. This is not a cause for concern.

E.367.2. Dump/Reload Procedure

If you are trying to reload a `pg_dump` or text-mode, `copy tablename` to stdout generated with a previous version, you will need to run the attached `sed` script on the ASCII file before loading it into the database. The old format used `.` as end-of-data, while `\` is now the end-of-data marker. Also, empty strings are now loaded in as `''` rather than NULL. See the copy manual page for full details.

```
sed 's/\.$/\./g' <in_file >out_file
```

If you are loading an older binary copy or non-stdout copy, there is no end-of-data character, and hence no conversion necessary.

```
-- following lines added by agc to reflect the case-insensitive
-- regexp searching for varchar (in 1.02), and bpchar (in 1.02.1)
create operator ~* (leftarg = bpchar, rightarg = text, procedure = texticregexeq);
create operator !~* (leftarg = bpchar, rightarg = text, procedure = texticregexne);
create operator ~* (leftarg = varchar, rightarg = text, procedure = texticregexeq);
create operator !~* (leftarg = varchar, rightarg = text, procedure = texticregexne);
```

E.367.3. Changes
Source code maintenance and development
* worldwide team of volunteers
* the source tree now in CVS at ftp.ki.net

Enhancements
* psql (and underlying libpq library) now has many more options for
  formatting output, including HTML
* pg_dump now output the schema and/or the data, with many fixes to
  enhance completeness.
* psql used in place of monitor in administration shell scripts.
  monitor to be deprecated in next release.
* date/time functions enhanced
* NULL insert/update/comparison fixed/enhanced
* TCL/TK lib and shell fixed to work with both tck7.4/tk4.0 and
tcl7.5/tk4.1

Bug Fixes (almost too numerous to mention)
* indexes
* storage management
* check for NULL pointer before dereferencing
* Makefile fixes

New Ports
* added SolarisX86 port
* added BSD/OS 2.1 port
* added DG/UX port

E.368. Release 1.01

Release date: 1996-02-23

E.368.1. Migration from version 1.0 to version 1.01

The following notes are for the benefit of users who want to migrate databases from Postgres95 1.0
to Postgres95 1.01.

If you are starting afresh with Postgres95 1.01 and do not need to migrate old databases, you do not
need to read any further.

In order to Postgres95 version 1.01 with databases created with Postgres95 version 1.0, the following
steps are required:
1. Set the definition of NAMEDATALEN in src/Makefile.global to 16 and OIDNAMELEN
to 20.
2. Decide whether you want to use Host based authentication.
   a. If you do, you must create a file name pg_hba in your top-level data directory (typically
      the value of your $PGDATA). src/libpq/pg_hba shows an example syntax.
   b. If you do not want host-based authentication, you can comment out the line:

      HBA = 1

      in src/Makefile.global
Note that host-based authentication is turned on by default, and if you do not take steps A or B above, the out-of-the-box 1.01 will not allow you to connect to 1.0 databases.

3. Compile and install 1.01, but DO NOT do the initdb step.

4. Before doing anything else, terminate your 1.0 postmaster, and backup your existing $PGDATA directory.

5. Set your $PGDATA environment variable to your 1.0 databases, but set up path up so that 1.01 binaries are being used.

6. Modify the file $PGDATA/PG_VERSION from 5.0 to 5.1.

7. Start up a new 1.01 postmaster.

8. Add the new built-in functions and operators of 1.01 to 1.0 databases. This is done by running the new 1.01 server against your own 1.0 database and applying the queries attached and saving in the file 1.0_to_1.01.sql. This can be done easily through psql. If your 1.0 database is name testdb:

% psql testdb -f 1.0_to_1.01.sql

and then execute the following commands (cut and paste from here):

```
-- add builtin functions that are new to 1.01
create function int4eqoid (int4, oid) returns bool as 'foo'
language 'internal';
create function oideqint4 (oid, int4) returns bool as 'foo'
language 'internal';
create function char2icregexeq (char2, text) returns bool as
'foo'
language 'internal';
create function char2icregexne (char2, text) returns bool as
'foo'
language 'internal';
create function char4icregexeq (char4, text) returns bool as
'foo'
language 'internal';
create function char4icregexne (char4, text) returns bool as
'foo'
language 'internal';
create function char8icregexeq (char8, text) returns bool as
'foo'
language 'internal';
create function char8icregexne (char8, text) returns bool as
'foo'
language 'internal';
create function char16icregexeq (char16, text) returns bool as
'foo'
language 'internal';
create function char16icregexne (char16, text) returns bool as
'foo'
language 'internal';
create function texticregexeq (text, text) returns bool as
'foo'
language 'internal';
```
create function texticregexne (text, text) returns bool as 'foo'
language 'internal';

-- add builtin functions that are new to 1.01

create operator = (leftarg = int4, rightarg = oid, procedure = int4eqoid);
create operator = (leftarg = oid, rightarg = int4, procedure = oideqint4);
create operator ~* (leftarg = char2, rightarg = text, procedure = char2icregexeq);
create operator !~* (leftarg = char2, rightarg = text, procedure = char2icregexne);
create operator ~* (leftarg = char4, rightarg = text, procedure = char4icregexeq);
create operator !~* (leftarg = char4, rightarg = text, procedure = char4icregexne);
create operator ~* (leftarg = char8, rightarg = text, procedure = char8icregexeq);
create operator !~* (leftarg = char8, rightarg = text, procedure = char8icregexne);
create operator ~* (leftarg = char16, rightarg = text, procedure = char16icregexeq);
create operator !~* (leftarg = char16, rightarg = text, procedure = char16icregexne);
create operator ~* (leftarg = text, rightarg = text, procedure = texticregexeq);
create operator !~* (leftarg = text, rightarg = text, procedure = texticregexne);

**E.368.2. Changes**

**Incompatibilities:**
* 1.01 is backwards compatible with 1.0 database provided the user
  follow the steps outlined in the MIGRATION_from_1.0_to_1.01
  file.
  If those steps are not taken, 1.01 is not compatible with 1.0
  database.

**Enhancements:**
* added PQdisplayTuples() to libpq and changed monitor and psql to
  use it
* added NeXT port (requires SysVIPC implementation)
* added CAST .. AS ... syntax
* added ASC and DESC key words
* added 'internal' as a possible language for CREATE FUNCTION
  internal functions are C functions which have been statically
  linked
  into the postgres backend.
* a new type "name" has been added for system identifiers (table
  names,
  attribute names, etc.) This replaces the old char16 type. The
  of name is set by the NAMEDATALEN #define in src/Makefile.global
* a readable reference manual that describes the query language.
* added host-based access control. A configuration file ($PGDATA/
  pg_hba)
is used to hold the configuration data. If host-based access control is not desired, comment out HBA=1 in src/Makefile.global.

* changed regex handling to be uniform use of Henry Spencer's regex code regardless of platform. The regex code is included in the distribution
* added functions and operators for case-insensitive regular expressions. The operators are ~* and !~*
* pg_dump uses COPY instead of SELECT loop for better performance

Bug fixes:
* fixed an optimizer bug that was causing core dumps when functions calls were used in comparisons in the WHERE clause
* changed all uses of getuid to geteuid so that effective uids are used
* psql now returns non-zero status on errors when using -c
* applied public patches 1-14

E.369. Release 1.0

Release date: 1995-09-05

E.369.1. Changes

Copyright change:
* The copyright of Postgres 1.0 has been loosened to be freely modifiable and modifiable for any purpose. Please read the COPYRIGHT file.
  Thanks to Professor Michael Stonebraker for making this possible.

Incompatibilities:
* date formats have to be MM-DD-YYYY (or DD-MM-YYYY if you're using EUROPEAN STYLE). This follows SQL-92 specs.
* "delimiters" is now a key word

Enhancements:
* sql LIKE syntax has been added
* copy command now takes an optional USING DELIMITER specification.
  delimiters can be any single-character string.
* IRIX 5.3 port has been added.
  Thanks to Paul Walmsley and others.
* updated pg_dump to work with new libpq
* \d has been added psql
  Thanks to Keith Parks
* regexp performance for architectures that use POSIX regex has been improved due to caching of precompiled patterns.
  Thanks to Alistair Crooks
* a new version of libpq++
  Thanks to William Wanders
Bug fixes:
* arbitrary userids can be specified in the createuser script
* \c to connect to other databases in psql now works.
* bad pg_proc entry for float4inc() is fixed
* users with usecreatedb field set can now create databases without
  having to be usesuper
* remove access control entries when the entry no longer has any
  privileges
* fixed non-portable datetimes implementation
* added kerberos flags to the src/backend/Makefile
* libpq now works with kerberos
* typographic errors in the user manual have been corrected.
* btrees with multiple index never worked, now we tell you they
  don't
  work when you try to use them

E.370. Postgres95 Release 0.03

Release date: 1995-07-21

E.370.1. Changes

Incompatible changes:
* BETA-0.3 IS INCOMPATIBLE WITH DATABASES CREATED WITH PREVIOUS
  VERSIONS
  (due to system catalog changes and indexing structure changes).
* double-quote (") is deprecated as a quoting character for string
  literals;
  you need to convert them to single quotes (').
* name of aggregates (eg. int4sum) are renamed in accordance with
  the
  SQL standard (eg. sum).
* CHANGE ACL syntax is replaced by GRANT/REVOKE syntax.
* float literals (eg. 3.14) are now of type float4 (instead of
  float8 in
  previous releases); you might have to do typecasting if you
  depend on it
  being of type float8. If you neglect to do the typecasting and
  you assign
  a float literal to a field of type float8, you might get
  incorrect values
  stored!
* LIBPQ has been totally revamped so that frontend applications
  can connect to multiple backends
* the usesysid field in pg_user has been changed from int2 to int4
  to
  allow wider range of Unix user ids.
* the netbsd/freebsd/bsd o/s ports have been consolidated into a
  single BSD44 Derived port. (thanks to Alistair Crooks)

SQL standard-compliance (the following details changes that makes
postgres95
more compliant to the SQL-92 standard):
* the following SQL types are now built-in: smallint, int(eger),
  float, real,
char(N), varchar(N), date and time.

The following are aliases to existing postgres types:
smallint -> int2
integer, int -> int4
float, real -> float4
char(N) and varchar(N) are implemented as truncated text types.
In addition, char(N) does blank-padding.
* single-quote (') is used for quoting string literals; '' (in addition to '\') is supported as means of inserting a single quote in a
string
* SQL standard aggregate names (MAX, MIN, AVG, SUM, COUNT) are used
(Also, aggregates can now be overloaded, i.e. you can define your
own MAX aggregate to take in a user-defined type.)
* CHANGE ACL removed. GRANT/REVOKE syntax added.
  - Privileges can be given to a group using the "GROUP" key word.
    For example:
    GRANT SELECT ON foobar TO GROUP my_group;
The key word 'PUBLIC' is also supported to mean all users.
Privileges can only be granted or revoked to one user or group
at a time.
"WITH GRANT OPTION" is not supported. Only class owners can change
access control
  - The default access control is to grant users readonly access.
You must explicitly grant insert/update access to users. To change
this, modify the line in
  src/backend/utils/acl.h
that defines ACL_WORLD_DEFAULT

Bug fixes:
* the bug where aggregates of empty tables were not run has been fixed. Now,
  aggregates run on empty tables will return the initial
  conditions of the aggregates. Thus, COUNT of an empty table will now properly
  return 0.
  MAX/MIN of an empty table will return a row of value NULL.
* allow the use of \; inside the monitor
* the LISTEN/NOTIFY asynchronous notification mechanism now work
* NOTIFY in rule action bodies now work
* hash indexes work, and access methods in general should perform better.
  creation of large btree indexes should be much faster. (thanks to Paul
  Aoki)

Other changes and enhancements:
* addition of an EXPLAIN statement used for explaining the query execution
plan (eg. "EXPLAIN SELECT * FROM EMP" prints out the execution plan for the query).
* WARN and NOTICE messages no longer have timestamps on them. To turn on timestamps of error messages, uncomment the line in src/backend/utils/elog.h:
  /* define ELOG_TIMESTAMP */
* On an access control violation, the message
  "Either no such class or insufficient privilege"
will be given. This is the same message that is returned when a class is not found. This dissuades non-privileged users from guessing the existence of privileged classes.
* some additional system catalog changes have been made that are not visible to the user.

libpqtcl changes:
* The -oid option has been added to the "pg_result" tcl command. pg_result -oid returns oid of the last row inserted. If the last command was not an INSERT, then pg_result -oid returns "."
* the large object interface is available as pg_lo* tcl commands: pg_lo_open, pg_lo_close, pg_lo_creat, etc.

Portability enhancements and New Ports:
* flex/lex problems have been cleared up. Now, you should be able to use flex instead of lex on any platforms. We no longer make assumptions of what lexer you use based on the platform you use.
* The Linux-ELF port is now supported. Various configuration have been tested: The following configuration is known to work:
  kernel 1.2.10, gcc 2.6.3, libc 4.7.2, flex 2.5.2, bison 1.24
  with everything in ELF format,

New utilities:
* ipcclean added to the distribution
  ipcclean usually does not need to be run, but if your backend crashes and leaves shared memory segments hanging around, ipcclean will clean them up for you.

New documentation:
* the user manual has been revised and libpq documentation added.

E.371. Postgres95 Release 0.02

Release date: 1995-05-25

E.371.1. Changes

Incompatible changes:
* The SQL statement for creating a database is 'CREATE DATABASE' instead
of 'CREATEDB'. Similarly, dropping a database is 'DROP DATABASE' instead of 'DESTROYDB'. However, the names of the executables 'createdb' and 'destroydb' remain the same.

New tools:
* pgperl - a Perl (4.036) interface to Postgres95
* pg_dump - a utility for dumping out a postgres database into a script file containing query commands. The script files are in an ASCII format and can be used to reconstruct the database, even on other machines and other architectures. (Also good for converting a Postgres 4.2 database to Postgres95 database.)

The following ports have been incorporated into postgres95-beta-0.02:
* the NetBSD port by Alistair Crooks
* the AIX port by Mike Tung
* the Windows NT port by Jon Forrest (more stuff but not done yet)
* the Linux ELF port by Brian Gallew

The following bugs have been fixed in postgres95-beta-0.02:
* new lines not escaped in COPY OUT and problem with COPY OUT when first attribute is a '.
* cannot type return to use the default user id in createuser
* SELECT DISTINCT on big tables crashes
* Linux installation problems
* monitor doesn't allow use of 'localhost' as PGHOST
* psql core dumps when doing \c or \l
* the "pgtclsh" target missing from src/bin/pgtclsh/Makefile
* libpgtcl has a hard-wired default port number
* SELECT DISTINCT INTO TABLE hangs
* CREATE TYPE doesn't accept 'variable' as the internal length
* wrong result using more than 1 aggregate in a SELECT

E.372. Postgres95 Release 0.01

Release date: 1995-05-01

Initial release.
Appendix F. Additional Supplied Modules

This appendix and the next one contain information regarding the modules that can be found in the contrib directory of the PostgreSQL distribution. These include porting tools, analysis utilities, and plug-in features that are not part of the core PostgreSQL system, mainly because they address a limited audience or are too experimental to be part of the main source tree. This does not preclude their usefulness.

This appendix covers extensions and other server plug-in modules found in contrib. Appendix G covers utility programs.

When building from the source distribution, these components are not built automatically, unless you build the "world" target (see Step 2). You can build and install all of them by running:

```
make
make install
```

in the contrib directory of a configured source tree; or to build and install just one selected module, do the same in that module's subdirectory. Many of the modules have regression tests, which can be executed by running:

```
make check
```

before installation or

```
make installcheck
```

once you have a PostgreSQL server running.

If you are using a pre-packaged version of PostgreSQL, these modules are typically made available as a separate subpackage, such as postgresql-contrib.

Many modules supply new user-defined functions, operators, or types. To make use of one of these modules, after you have installed the code you need to register the new SQL objects in the database system. In PostgreSQL 9.1 and later, this is done by executing a CREATE EXTENSION command. In a fresh database, you can simply do

```
CREATE EXTENSION module_name;
```

This command must be run by a database superuser. This registers the new SQL objects in the current database only, so you need to run this command in each database that you want the module's facilities to be available in. Alternatively, run it in database template1 so that the extension will be copied into subsequently-created databases by default.

Many modules allow you to install their objects in a schema of your choice. To do that, add SCHEMA schema_name to the CREATE EXTENSION command. By default, the objects will be placed in your current creation target schema, which in turn defaults to public.

If your database was brought forward by dump and reload from a pre-9.1 version of PostgreSQL, and you had been using the pre-9.1 version of the module in it, you should instead do
CREATE EXTENSION module_name FROM unpackaged;

This will update the pre-9.1 objects of the module into a proper extension object. Future updates to the module will be managed by ALTER EXTENSION. For more information about extension updates, see Section 38.16.

Note, however, that some of these modules are not "extensions" in this sense, but are loaded into the server in some other way, for instance by way of shared_preload_libraries. See the documentation of each module for details.

**F.1. adminpack**

adminpack provides a number of support functions which pgAdmin and other administration and management tools can use to provide additional functionality, such as remote management of server log files. Use of all these functions is only allowed to the superuser by default but may be allowed to other users by using the GRANT command.

The functions shown in Table F.1 provide write access to files on the machine hosting the server. (See also the functions in Table 9.88, which provide read-only access.) Only files within the database cluster directory can be accessed, unless the user is a superuser or given one of the pg_read_server_files, or pg_write_server_files roles, as appropriate for the function, but either a relative or absolute path is allowable.

**Table F.1. adminpack Functions**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_cata-log.pg_file_write(filename text, data text, append boolean)</td>
<td>bigint</td>
<td>Write, or append to, a text file</td>
</tr>
<tr>
<td>pg_cata-log.pg_file_rename(oldname text, newname text [, archivename text])</td>
<td>boolean</td>
<td>Rename a file</td>
</tr>
<tr>
<td>pg_cata-log.pg_file_unlink(filename text)</td>
<td>boolean</td>
<td>Remove a file</td>
</tr>
<tr>
<td>pg_cata-log.pg_logdir_ls()</td>
<td>setof record</td>
<td>List the log files in the log_directory directory</td>
</tr>
</tbody>
</table>

pg_file_write writes the specified data into the file named by filename. If append is false, the file must not already exist. If append is true, the file can already exist, and will be appended to if so. Returns the number of bytes written.

pg_file_rename renames a file. If archivename is omitted or NULL, it simply renames oldname to newname (which must not already exist). If archivename is provided, it first renames newname to archivename (which must not already exist), and then renames oldname to newname. In event of failure of the second rename step, it will try to rename archivename back to newname before reporting the error. Returns true on success, false if the source file(s) are not present or not writable; other cases throw errors.

pg_file_unlink removes the specified file. Returns true on success, false if the specified file is not present or the unlink() call fails; other cases throw errors.
pg_logdir_ls returns the start timestamps and path names of all the log files in the log_directory directory. The log_filename parameter must have its default setting (postgresql-%Y-%m-%d-%H%M%S.log) to use this function.

F.2. amcheck

The amcheck module provides functions that allow you to verify the logical consistency of the structure of relations. If the structure appears to be valid, no error is raised.

The functions verify various invariants in the structure of the representation of particular relations. The correctness of the access method functions behind index scans and other important operations relies on these invariants always holding. For example, certain functions verify, among other things, that all B-Tree pages have items in “logical” order (e.g., for B-Tree indexes on text, index tuples should be in collated lexical order). If that particular invariant somehow fails to hold, we can expect binary searches on the affected page to incorrectly guide index scans, resulting in wrong answers to SQL queries.

Verification is performed using the same procedures as those used by index scans themselves, which may be user-defined operator class code. For example, B-Tree index verification relies on comparisons made with one or more B-Tree support function 1 routines. See Section 38.15.3 for details of operator class support functions.

amcheck functions may only be used by superusers.

F.2.1. Functions

bt_index_check(index regclass, heapallindexed boolean) returns void

bt_index_check tests that its target, a B-Tree index, respects a variety of invariants. Example usage:

test=# SELECT bt_index_check(index => c.oid, heapallindexed => i.indisunique),
       c.relname,
       c.relpages
FROM pg_index i
JOIN pg_opclass op ON i.indclass[0] = op.oid
JOIN pg_am am ON op.opcmethod = am.oid
JOIN pg_class c ON i.indexrelid = c.oid
JOIN pg_namespace n ON c.relnamespace = n.oid
WHERE am.amname = 'btree' AND n.nspname = 'pg_catalog'
   -- Don't check temp tables, which may be from another session:
   AND c.relpersistence != 't'
   -- Function may throw an error when this is omitted:
   AND c.relkind = 'i' AND i.indisready AND i.indisvalid
ORDER BY c.relpages DESC LIMIT 10;

<table>
<thead>
<tr>
<th>bt_index_check</th>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_depend_reference_index</td>
<td>pg_depender_index</td>
<td>43</td>
</tr>
<tr>
<td>pg_depender_index</td>
<td>pgproc_proname_args_nsp_index</td>
<td>40</td>
</tr>
<tr>
<td>pgproc_proname_args_nsp_index</td>
<td>pg_description_o_c_o_index</td>
<td>31</td>
</tr>
<tr>
<td>pg_description_o_c_o_index</td>
<td>pg_attribute_relid_attnam_index</td>
<td>21</td>
</tr>
<tr>
<td>pg_attribute_relid_attnam_index</td>
<td>pgproc_oid_index</td>
<td>14</td>
</tr>
<tr>
<td>pgproc_oid_index</td>
<td>pg_attribute_relid_attnum_index</td>
<td>10</td>
</tr>
<tr>
<td>pg_attribute_relid_attnum_index</td>
<td>pgamproc_fam_proc_index</td>
<td>9</td>
</tr>
<tr>
<td>pgamproc_fam_proc_index</td>
<td>pgamop_opr_fam_index</td>
<td>5</td>
</tr>
</tbody>
</table>
This example shows a session that performs verification of the 10 largest catalog indexes in the database “test”. Verification of the presence of heap tuples as index tuples is requested for the subset that are unique indexes. Since no error is raised, all indexes tested appear to be logically consistent. Naturally, this query could easily be changed to call bt_index_check for every index in the database where verification is supported.

bt_index_check acquires an AccessShareLock on the target index and the heap relation it belongs to. This lock mode is the same lock mode acquired on relations by simple SELECT statements. bt_index_check does not verify invariants that span child/parent relationships, but will verify the presence of all heap tuples as index tuples within the index when heapallindexed is true. When a routine, lightweight test for corruption is required in a live production environment, using bt_index_check often provides the best trade-off between thoroughness of verification and limiting the impact on application performance and availability.

bt_index_parent_check(index regclass, heapallindexed boolean) returns void

bt_index_parent_check tests that its target, a B-Tree index, respects a variety of invariants. Optionally, when the heapallindexed argument is true, the function verifies the presence of all heap tuples that should be found within the index, and that there are no missing down-links in the index structure. The checks that can be performed by bt_index_parent_check are a superset of the checks that can be performed by bt_index_check. bt_index_parent_check can be thought of as a more thorough variant of bt_index_check: unlike bt_index_check, bt_index_parent_check also checks invariants that span parent/child relationships. bt_index_parent_check follows the general convention of raising an error if it finds a logical inconsistency or other problem.

A ShareLock is required on the target index by bt_index_parent_check (a ShareLock is also acquired on the heap relation). These locks prevent concurrent data modification from INSERT, UPDATE, and DELETE commands. The locks also prevent the underlying relation from being concurrently processed by VACUUM, as well as all other utility commands. Note that the function holds locks only while running, not for the entire transaction.

bt_index_parent_check’s additional verification is more likely to detect various pathological cases. These cases may involve an incorrectly implemented B-Tree operator class used by the index that is checked, or, hypothetically, undiscovered bugs in the underlying B-Tree index access method code. Note that bt_index_parent_check cannot be used when Hot Standby mode is enabled (i.e., on read-only physical replicas), unlike bt_index_check.

F.2.2. Optional heapallindexed verification

When the heapallindexed argument to verification functions is true, an additional phase of verification is performed against the table associated with the target index relation. This consists of a “dummy” CREATE INDEX operation, which checks for the presence of all hypothetical new index tuples against a temporary, in-memory summarizing structure (this is built when needed during the basic first phase of verification). The summarizing structure “fingerprint” every tuple found within the target index. The high level principle behind heapallindexed verification is that a new index that is equivalent to the existing, target index must only have entries that can be found in the existing structure.

The additional heapallindexed phase adds significant overhead: verification will typically take several times longer. However, there is no change to the relation-level locks acquired when heapallindexed verification is performed.

The summarizing structure is bound in size by maintenance_work_mem. In order to ensure that there is no more than a 2% probability of failure to detect an inconsistency for each heap tuple that should be represented in the index, approximately 2 bytes of memory are needed per tuple. As less
memory is made available per tuple, the probability of missing an inconsistency slowly increases. This approach limits the overhead of verification significantly, while only slightly reducing the probability of detecting a problem, especially for installations where verification is treated as a routine maintenance task. Any single absent or malformed tuple has a new opportunity to be detected with each new verification attempt.

F.2.3. Using amcheck effectively

amcheck can be effective at detecting various types of failure modes that data page checksums will always fail to catch. These include:

- Structural inconsistencies caused by incorrect operator class implementations.

  This includes issues caused by the comparison rules of operating system collations changing. Comparisons of datums of a collatable type like `text` must be immutable (just as all comparisons used for B-Tree index scans must be immutable), which implies that operating system collation rules must never change. Though rare, updates to operating system collation rules can cause these issues. More commonly, an inconsistency in the collation order between a master server and a standby server is implicated, possibly because the major operating system version in use is inconsistent. Such inconsistencies will generally only arise on standby servers, and so can generally only be detected on standby servers.

  If a problem like this arises, it may not affect each individual index that is ordered using an affected collation, simply because indexed values might happen to have the same absolute ordering regardless of the behavioral inconsistency. See Section 23.1 and Section 23.2 for further details about how PostgreSQL uses operating system locales and collations.

- Structural inconsistencies between indexes and the heap relations that are indexed (when `heapallindexed` verification is performed).

  There is no cross-checking of indexes against their heap relation during normal operation. Symptoms of heap corruption can be subtle.

- Corruption caused by hypothetical undiscovered bugs in the underlying PostgreSQL access method code, sort code, or transaction management code.

  Automatic verification of the structural integrity of indexes plays a role in the general testing of new or proposed PostgreSQL features that could plausibly allow a logical inconsistency to be introduced. Verification of table structure and associated visibility and transaction status information plays a similar role. One obvious testing strategy is to call `amcheck` functions continuously when running the standard regression tests. See Section 33.1 for details on running the tests.

- File system or storage subsystem faults where checksums happen to simply not be enabled.

  Note that `amcheck` examines a page as represented in some shared memory buffer at the time of verification if there is only a shared buffer hit when accessing the block. Consequently, `amcheck` does not necessarily examine data read from the file system at the time of verification. Note that when checksums are enabled, `amcheck` may raise an error due to a checksum failure when a corrupt block is read into a buffer.

- Corruption caused by faulty RAM, or the broader memory subsystem.

  PostgreSQL does not protect against correctable memory errors and it is assumed you will operate using RAM that uses industry standard Error Correcting Codes (ECC) or better protection. However, ECC memory is typically only immune to single-bit errors, and should not be assumed to provide absolute protection against failures that result in memory corruption.

  When `heapallindexed` verification is performed, there is generally a greatly increased chance of detecting single-bit errors, since strict binary equality is tested, and the indexed attributes within the heap are tested.
In general, `amcheck` can only prove the presence of corruption; it cannot prove its absence.

**F.2.4. Repairing corruption**

No error concerning corruption raised by `amcheck` should ever be a false positive. `amcheck` raises errors in the event of conditions that, by definition, should never happen, and so careful analysis of `amcheck` errors is often required.

There is no general method of repairing problems that `amcheck` detects. An explanation for the root cause of an invariant violation should be sought. `pageinspect` may play a useful role in diagnosing corruption that `amcheck` detects. A `REINDEX` may not be effective in repairing corruption.

**F.3. auth_delay**

`auth_delay` causes the server to pause briefly before reporting authentication failure, to make brute-force attacks on database passwords more difficult. Note that it does nothing to prevent denial-of-service attacks, and may even exacerbate them, since processes that are waiting before reporting authentication failure will still consume connection slots.

In order to function, this module must be loaded via `shared_preload_libraries` in `postgresql.conf`.

**F.3.1. Configuration Parameters**

`auth_delay.milliseconds (int)`

The number of milliseconds to wait before reporting an authentication failure. The default is 0.

These parameters must be set in `postgresql.conf`. Typical usage might be:

```bash
# postgresql.conf
shared_preload_libraries = 'auth_delay'
auth_delay.milliseconds = '500'
```

**F.3.2. Author**

KaiGai Kohei <kaigai@ak.jp.nec.com>

**F.4. auto_explain**

The `auto_explain` module provides a means for logging execution plans of slow statements automatically, without having to run `EXPLAIN` by hand. This is especially helpful for tracking down un-optimized queries in large applications.

The module provides no SQL-accessible functions. To use it, simply load it into the server. You can load it into an individual session:

```sql
LOAD 'auto_explain';
```

(You must be superuser to do that.) More typical usage is to preload it into some or all sessions by including `auto_explain` in `session_preload_libraries` or `shared_preload_libraries` in `postgresql.conf`. Then you can track unexpectedly slow queries no matter when they happen. Of course there is a price in overhead for that.
Additional Supplied Modules

F.4.1. Configuration Parameters

There are several configuration parameters that control the behavior of `auto_explain`. Note that the default behavior is to do nothing, so you must set at least `auto_explain.log_min_duration` if you want any results.

`auto_explain.log_min_duration` (integer)

`auto_explain.log_min_duration` is the minimum statement execution time, in milliseconds, that will cause the statement's plan to be logged. Setting this to zero logs all plans. Minus-one (the default) disables logging of plans. For example, if you set it to 250ms then all statements that run 250ms or longer will be logged. Only superusers can change this setting.

`auto_explain.log_analyze` (boolean)

`auto_explain.log_analyze` causes `EXPLAIN ANALYZE` output, rather than just `EXPLAIN` output, to be printed when an execution plan is logged. This parameter is off by default. Only superusers can change this setting.

**Note**

When this parameter is on, per-plan-node timing occurs for all statements executed, whether or not they run long enough to actually get logged. This can have an extremely negative impact on performance. Turning off `auto_explain.log_timing` ameliorates the performance cost, at the price of obtaining less information.

`auto_explain.log_buffers` (boolean)

`auto_explain.log_buffers` controls whether buffer usage statistics are printed when an execution plan is logged; it's equivalent to the `BUFFERS` option of `EXPLAIN`. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is off by default. Only superusers can change this setting.

`auto_explain.log_timing` (boolean)

`auto_explain.log_timing` controls whether per-node timing information is printed when an execution plan is logged; it's equivalent to the `TIMING` option of `EXPLAIN`. The overhead of repeatedly reading the system clock can slow down queries significantly on some systems, so it may be useful to set this parameter to off when only actual row counts, and not exact times, are needed. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is on by default. Only superusers can change this setting.

`auto_explain.log_triggers` (boolean)

`auto_explain.log_triggers` causes trigger execution statistics to be included when an execution plan is logged. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is off by default. Only superusers can change this setting.

`auto_explain.log_verbose` (boolean)

`auto_explain.log_verbose` controls whether verbose details are printed when an execution plan is logged; it's equivalent to the `VERBOSE` option of `EXPLAIN`. This parameter is off by default. Only superusers can change this setting.

`auto_explain.log_format` (enum)

`auto_explain.log_format` selects the `EXPLAIN` output format to be used. The allowed values are `text`, `xml`, `json`, and `yaml`. The default is `text`. Only superusers can change this setting.
**Additional Supplied Modules**

**auto_explain.log_nested_statements** *(boolean)*

`auto_explain.log_nested_statements` causes nested statements (statements executed inside a function) to be considered for logging. When it is off, only top-level query plans are logged. This parameter is off by default. Only superusers can change this setting.

**auto_explain.sample_rate** *(real)*

`auto_explain.sample_rate` causes `auto_explain` to only explain a fraction of the statements in each session. The default is 1, meaning explain all the queries. In case of nested statements, either all will be explained or none. Only superusers can change this setting.

In ordinary usage, these parameters are set in `postgresql.conf`, although superusers can alter them on-the-fly within their own sessions. Typical usage might be:

```plaintext
# postgresql.conf
session_preload_libraries = 'auto_explain'
auto_explain.log_min_duration = '3s'
```

**F.4.2. Example**

```sql
postgres=# LOAD 'auto_explain';
p没见过# SET auto_explain.log_min_duration = 0;
p没见过# SET auto_explain.log_analyze = true;
p没见过# SELECT count(*)
FROM pg_class, pg_index
WHERE oid = indrelid AND indisunique;
```

This might produce log output such as:

```
LOG: duration: 3.651 ms  plan:
  Query Text: SELECT count(*)
               FROM pg_class, pg_index
               WHERE oid = indrelid AND indisunique;
  Aggregate (cost=16.79..16.80 rows=1 width=0) (actual time=3.626 ..3.627 rows=1 loops=1)
    -> Hash Join (cost=4.17 ..16.55 rows=92 width=0) (actual time=3.349 ..3.594 rows=92 loops=1)
      Hash Cond: (pg_class.oid = pg_index.indrelid)
      -> Seq Scan on pg_class (cost=0.00 ..9.55 rows=255 width=4) (actual time=0.016 ..0.140 rows=255 loops=1)
      -> Hash (cost=3.02 ..3.02 rows=92 width=4) (actual time=3.238 ..3.238 rows=92 loops=1)
        Buckets: 1024  Batches: 1  Memory Usage: 4kB
        -> Seq Scan on pg_index (cost=0.00 ..3.02 rows=92 width=4) (actual time=0.008 ..3.187 rows=92 loops=1)
          Filter: indisunique
```

**F.4.3. Author**

Takahiro Itagaki <itagaki.takahiro@oss.ntt.co.jp>

**F.5. bloom**
bloom provides an index access method based on Bloom filters\(^1\).

A Bloom filter is a space-efficient data structure that is used to test whether an element is a member of a set. In the case of an index access method, it allows fast exclusion of non-matching tuples via signatures whose size is determined at index creation.

A signature is a lossy representation of the indexed attribute(s), and as such is prone to reporting false positives; that is, it may be reported that an element is in the set, when it is not. So index search results must always be rechecked using the actual attribute values from the heap entry. Larger signatures reduce the odds of a false positive and thus reduce the number of useless heap visits, but of course also make the index larger and hence slower to scan.

This type of index is most useful when a table has many attributes and queries test arbitrary combinations of them. A traditional btree index is faster than a bloom index, but it can require many btree indexes to support all possible queries where one needs only a single bloom index. Note however that bloom indexes only support equality queries, whereas btree indexes can also perform inequality and range searches.

**F.5.1. Parameters**

A bloom index accepts the following parameters in its `WITH` clause:

- **length**
  
  Length of each signature (index entry) in bits. It is rounded up to the nearest multiple of 16. The default is 80 bits and the maximum is 4096.

- **col1** — **col32**
  
  Number of bits generated for each index column. Each parameter's name refers to the number of the index column that it controls. The default is 2 bits and maximum is 4095. Parameters for index columns not actually used are ignored.

**F.5.2. Examples**

This is an example of creating a bloom index:

```sql
CREATE INDEX bloomidx ON tbloom USING bloom (i1,i2,i3)
  WITH (length=80, col1=2, col2=2, col3=4);
```

The index is created with a signature length of 80 bits, with attributes i1 and i2 mapped to 2 bits, and attribute i3 mapped to 4 bits. We could have omitted the `length`, `col1`, and `col2` specifications since those have the default values.

Here is a more complete example of bloom index definition and usage, as well as a comparison with equivalent btree indexes. The bloom index is considerably smaller than the btree index, and can perform better.

```sql
=# CREATE TABLE tbloom AS
  SELECT
    (random() * 1000000)::int as i1,
    (random() * 1000000)::int as i2,
    (random() * 1000000)::int as i3,
    (random() * 1000000)::int as i4,
    (random() * 1000000)::int as i5,
    (random() * 1000000)::int as i6
```

\(^1\) [https://en.wikipedia.org/wiki/Bloom_filter](https://en.wikipedia.org/wiki/Bloom_filter)
FROM generate_series(1,10000000);
SELECT 10000000
=# CREATE INDEX bloomidx ON tbloom USING bloom (i1, i2, i3, i4, i5, i6);
CREATE INDEX
=# SELECT pg_size_pretty(pg_relation_size('bloomidx'));
pg_size_pretty
----------------
153 MB
(1 row)
=# CREATE index btreeidx ON tbloom (i1, i2, i3, i4, i5, i6);
CREATE INDEX
=# SELECT pg_size_pretty(pg_relation_size('btreeidx'));
pg_size_pretty
----------------
387 MB
(1 row)

A sequential scan over this large table takes a long time:

=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
QUERY PLAN
--------------------------------------------------------------------------------------------------------------------------------
Seq Scan on tbloom (cost=0.00..213694.08 rows=1 width=24) (actual time=1445.438..1445.438 rows=0 loops=1)
  Filter: ((i2 = 898732) AND (i5 = 123451))
Rows Removed by Filter: 10000000
Planning time: 0.177 ms
Execution time: 1445.473 ms
(5 rows)

So the planner will usually select an index scan if possible. With a btree index, we get results like this:

=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
QUERY PLAN
---------------------------------------------------------------------------------------------------------------------------
Bitmap Heap Scan on tbloom (cost=178435.39..178439.41 rows=1 width=24) (actual time=76.698..76.698 rows=0 loops=1)

Bloom is better than btree in handling this type of search:
Recheck Cond: ((i2 = 898732) AND (i5 = 123451))
Rows Removed by Index Recheck: 2439
Heap Blocks: exact=2408
-> Bitmap Index Scan on bloomidx (cost=0.00..178435.39 rows=1 width=0) (actual time=72.455..72.455 rows=2439 loops=1)
  Index Cond: ((i2 = 898732) AND (i5 = 123451))
Planning time: 0.475 ms
Execution time: 76.778 ms
(8 rows)

Note the relatively large number of false positives: 2439 rows were selected to be visited in the heap, but none actually matched the query. We could reduce that by specifying a larger signature length. In this example, creating the index with length=200 reduced the number of false positives to 55; but it doubled the index size (to 306 MB) and ended up being slower for this query (125 ms overall).

Now, the main problem with the btree search is that btree is inefficient when the search conditions do not constrain the leading index column(s). A better strategy for btree is to create a separate index on each column. Then the planner will choose something like this:

```sql
=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
```

```sql
PLAN

Bitmap Heap Scan on tbloom (cost=9.29..13.30 rows=1 width=24)
  (actual time=0.148..0.148 rows=0 loops=1)
  Recheck Cond: ((i5 = 123451) AND (i2 = 898732))
  -> BitmapAnd (cost=9.29..9.29 rows=1 width=0) (actual time=0.145..0.145 rows=0 loops=1)
    -> Bitmap Index Scan on tbloom_i5_idx (cost=0.00..4.52 rows=11 width=0) (actual time=0.089..0.089 rows=10 loops=1)
      Index Cond: (i5 = 123451)
    -> Bitmap Index Scan on tbloom_i2_idx (cost=0.00..4.52 rows=11 width=0) (actual time=0.048..0.048 rows=8 loops=1)
      Index Cond: (i2 = 898732)
Planning time: 2.049 ms
Execution time: 0.280 ms
(9 rows)
```

Although this query runs much faster than with either of the single indexes, we pay a large penalty in index size. Each of the single-column btree indexes occupies 214 MB, so the total space needed is over 1.2GB, more than 8 times the space used by the bloom index.

### F.5.3. Operator Class Interface

An operator class for bloom indexes requires only a hash function for the indexed data type and an equality operator for searching. This example shows the operator class definition for the `text` data type:

```sql
CREATE OPERATOR CLASS text_ops
DEFAULT FOR TYPE text USING bloom AS
  OPERATOR 1 =(text, text),
  FUNCTION 1 hashtext(text);
```

### F.5.4. Limitations

- Only operator classes for `int4` and `text` are included with the module.
• Only the = operator is supported for search. But it is possible to add support for arrays with union and intersection operations in the future.

• bloom access method doesn’t support UNIQUE indexes.

• bloom access method doesn’t support searching for NULL values.

F.5.5. Authors

Teodor Sigaev <teodor@postgrespro.ru>, Postgres Professional, Moscow, Russia

Alexander Korotkov <a.korotkov@postgrespro.ru>, Postgres Professional, Moscow, Russia

Oleg Bartunov <obartunov@postgrespro.ru>, Postgres Professional, Moscow, Russia

F.6. btrees_gin

btrees_gin provides sample GIN operator classes that implement B-tree equivalent behavior for the data types int2, int4, int8, float4, float8, timestamp with time zone, timestamp without time zone, time with time zone, time without time zone, date, interval, oid, money, "char", varchar, text, bytea, bit, varbit, macaddr, macaddr8, inet, cidr, uuid, name, bool, bpchar, and all enum types.

In general, these operator classes will not outperform the equivalent standard B-tree index methods, and they lack one major feature of the standard B-tree code: the ability to enforce uniqueness. However, they are useful for GIN testing and as a base for developing other GIN operator classes. Also, for queries that test both a GIN-indexable column and a B-tree-indexable column, it might be more efficient to create a multicolumn GIN index that uses one of these operator classes than to create two separate indexes that would have to be combined via bitmap ANDing.

F.6.1. Example Usage

```
CREATE TABLE test (a int4);
-- create index
CREATE INDEX testidx ON test USING GIN (a);
-- query
SELECT * FROM test WHERE a < 10;
```

F.6.2. Authors

Teodor Sigaev (<teodor@stack.net>) and Oleg Bartunov (<oleg@sai.msu.su>). See http://www.sai.msu.su/~megera/oddmuse/index.cgi/Gin for additional information.

F.7. btrees_gist

btrees_gist provides GiST index operator classes that implement B-tree equivalent behavior for the data types int2, int4, int8, float4, float8, numeric, timestamp with time zone, timestamp without time zone, time with time zone, time without time zone, date, interval, oid, money, char, varchar, text, bytea, bit, varbit, macaddr, macaddr8, inet, cidr, uuid, and all enum types.

In general, these operator classes will not outperform the equivalent standard B-tree index methods, and they lack one major feature of the standard B-tree code: the ability to enforce uniqueness. However, they provide some other features that are not available with a B-tree index, as described below. Also, these operator classes are useful when a multicolumn GiST index is needed, wherein some of the
columns are of data types that are only indexable with GiST but other columns are just simple data types. Lastly, these operator classes are useful for GiST testing and as a base for developing other GiST operator classes.

In addition to the typical B-tree search operators, btree_gist also provides index support for <> ("not equals"). This may be useful in combination with an exclusion constraint, as described below.

Also, for data types for which there is a natural distance metric, btree_gist defines a distance operator <->, and provides GiST index support for nearest-neighbor searches using this operator. Distance operators are provided for int2, int4, int8, float4, float8, timestamp with time zone, timestamp without time zone, time without time zone, date, interval, oid, and money.

F.7.1. Example Usage

Simple example using btree_gist instead of btree:

```sql
CREATE TABLE test (a int4);
-- create index
CREATE INDEX testidx ON test USING GIST (a);
-- query
SELECT * FROM test WHERE a < 10;
-- nearest-neighbor search: find the ten entries closest to "42"
SELECT *, a <-> 42 AS dist FROM test ORDER BY a <-> 42 LIMIT 10;
```

Use an exclusion constraint to enforce the rule that a cage at a zoo can contain only one kind of animal:

```sql
=> CREATE TABLE zoo (  
cage INTEGER,  
animal TEXT,  
EXCLUDE USING GIST (cage WITH =, animal WITH <>)
);
=> INSERT INTO zoo VALUES(123, 'zebra');
INSERT 0 1
=> INSERT INTO zoo VALUES(123, 'zebra');
INSERT 0 1
=> INSERT INTO zoo VALUES(123, 'lion');
ERROR: conflicting key value violates exclusion constraint
"zoo_cage_animal_excl"
DETAIL: Key (cage, animal)=(123, lion) conflicts with existing key
(cage, animal)=(123, zebra).
=> INSERT INTO zoo VALUES(124, 'lion');
INSERT 0 1
```

F.7.2. Authors

Teodor Sigaev (<teodor@stack.net>), Oleg Bartunov (<oleg@sai.msu.su>), Janko Richter (<jankorichter@yahoo.de>), and Paul Jungwirth (<pj@illuminatedcomputing.com>). See http://www.sai.msu.su/~megera/postgres/gist/ for additional information.

F.8. citext

The citext module provides a case-insensitive character string type, citext. Essentially, it internally calls lower when comparing values. Otherwise, it behaves almost exactly like text.
F.8.1. Rationale

The standard approach to doing case-insensitive matches in PostgreSQL has been to use the `lower` function when comparing values, for example:

```
SELECT * FROM tab WHERE lower(col) = LOWER(?);
```

This works reasonably well, but has a number of drawbacks:

- It makes your SQL statements verbose, and you always have to remember to use `lower` on both the column and the query value.
- It won't use an index, unless you create a functional index using `lower`.
- If you declare a column as `UNIQUE` or `PRIMARY KEY`, the implicitly generated index is case-sensitive. So it's useless for case-insensitive searches, and it won't enforce uniqueness case-insensitively.

The `citext` data type allows you to eliminate calls to `lower` in SQL queries, and allows a primary key to be case-insensitive. `citext` is locale-aware, just like `text`, which means that the matching of upper case and lower case characters is dependent on the rules of the database's `LC_CTYPE` setting. Again, this behavior is identical to the use of `lower` in queries. But because it's done transparently by the data type, you don't have to remember to do anything special in your queries.

F.8.2. How to Use It

Here's a simple example of usage:

```
CREATE TABLE users (
    nick CITEXT PRIMARY KEY,
    pass TEXT NOT NULL
);

INSERT INTO users VALUES ( 'larry', sha256(random()::text::bytea) );
INSERT INTO users VALUES ( 'Tom', sha256(random()::text::bytea) );
INSERT INTO users VALUES ( 'Damian', sha256(random()::text::bytea) );
INSERT INTO users VALUES ( 'NEAL', sha256(random()::text::bytea) );
INSERT INTO users VALUES ( 'Bjørn', sha256(random()::text::bytea) );

SELECT * FROM users WHERE nick = 'Larry';
```

The `SELECT` statement will return one tuple, even though the `nick` column was set to `larry` and the query was for `Larry`.

F.8.3. String Comparison Behavior

`citext` performs comparisons by converting each string to lower case (as though `lower` were called) and then comparing the results normally. Thus, for example, two strings are considered equal if `lower` would produce identical results for them.

In order to emulate a case-insensitive collation as closely as possible, there are `citext`-specific versions of a number of string-processing operators and functions. So, for example, the regular expres-
sion operators ~ and ~* exhibit the same behavior when applied to citext: they both match case-insensitively. The same is true for !~ and !~*, as well as for the LIKE operators ~~ and ~*, and !~~ and !~*. If you'd like to match case-sensitively, you can cast the operator's arguments to text.

Similarly, all of the following functions perform matching case-insensitively if their arguments are citext:

- `regexp_match()`
- `regexp_matches()`
- `regexp_replace()`
- `regexp_split_to_array()`
- `regexp_split_to_table()`
- `replace()`
- `split_part()`
- `strpos()`
- `translate()`

For the regexp functions, if you want to match case-sensitively, you can specify the "c" flag to force a case-sensitive match. Otherwise, you must cast to text before using one of these functions if you want case-sensitive behavior.

### F.8.4. Limitations

- citext's case-folding behavior depends on the LC_CTYPE setting of your database. How it compares values is therefore determined when the database is created. It is not truly case-insensitive in the terms defined by the Unicode standard. Effectively, what this means is that, as long as you're happy with your collation, you should be happy with citext's comparisons. But if you have data in different languages stored in your database, users of one language may find their query results are not as expected if the collation is for another language.

- As of PostgreSQL 9.1, you can attach a COLLATE specification to citext columns or data values. Currently, citext operators will honor a non-default COLLATE specification while comparing case-folded strings, but the initial folding to lower case is always done according to the database's LC_CTYPE setting (that is, as though COLLATE "default" were given). This may be changed in a future release so that both steps follow the input COLLATE specification.

- citext is not as efficient as text because the operator functions and the B-tree comparison functions must make copies of the data and convert it to lower case for comparisons. It is, however, slightly more efficient than using lower to get case-insensitive matching.

- citext doesn't help much if you need data to compare case-sensitively in some contexts and case-insensitively in other contexts. The standard answer is to use the text type and manually use the lower function when you need to compare case-insensitively; this works all right if case-insensitive comparison is needed only infrequently. If you need case-insensitive behavior most of the time and case-sensitive infrequently, consider storing the data as citext and explicitly casting the column to text when you want case-sensitive comparison. In either situation, you will need two indexes if you want both types of searches to be fast.

- The schema containing the citext operators must be in the current search_path (typically public); if it is not, the normal case-sensitive text operators will be invoked instead.

### F.8.5. Author

David E. Wheeler <david@kineticode.com>
Additional Supplied Modules

Inspired by the original citext module by Donald Fraser.

F.9. cube

This module implements a data type cube for representing multidimensional cubes.

F.9.1. Syntax

Table F.2 shows the valid external representations for the cube type. x, y, etc. denote floating-point numbers.

Table F.2. Cube External Representations

<table>
<thead>
<tr>
<th>External Syntax</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>x</td>
<td>A one-dimensional point (or, zero-length one-dimensional interval)</td>
</tr>
<tr>
<td>(x)</td>
<td>Same as above</td>
</tr>
<tr>
<td>x1,x2,...,xn</td>
<td>A point in n-dimensional space, represented internally as a zero-volume cube</td>
</tr>
<tr>
<td>(x1,x2,...,xn)</td>
<td>Same as above</td>
</tr>
<tr>
<td>(x), (y)</td>
<td>A one-dimensional interval starting at x and ending at y or vice versa; the order does not matter</td>
</tr>
<tr>
<td>[(x), (y)]</td>
<td>Same as above</td>
</tr>
<tr>
<td>(x1,...,xn), (y1,...,yn)</td>
<td>An n-dimensional cube represented by a pair of its diagonally opposite corners</td>
</tr>
<tr>
<td>[(x1,...,xn), (y1,...,yn)]</td>
<td>Same as above</td>
</tr>
</tbody>
</table>

It does not matter which order the opposite corners of a cube are entered in. The cube functions automatically swap values if needed to create a uniform “lower left — upper right” internal representation. When the corners coincide, cube stores only one corner along with an “is point” flag to avoid wasting space.

White space is ignored on input, so [(x), (y)] is the same as [(x), (y)].

F.9.2. Precision

Values are stored internally as 64-bit floating point numbers. This means that numbers with more than about 16 significant digits will be truncated.

F.9.3. Usage

Table F.3 shows the operators provided for type cube.

Table F.3. Cube Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Result</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>a = b</td>
<td>boolean</td>
<td>The cubes a and b are identical.</td>
</tr>
<tr>
<td>a &amp;&amp; b</td>
<td>boolean</td>
<td>The cubes a and b overlap.</td>
</tr>
<tr>
<td>a @&gt; b</td>
<td>boolean</td>
<td>The cube a contains the cube b.</td>
</tr>
<tr>
<td>a &lt;@ b</td>
<td>boolean</td>
<td>The cube a is contained in the cube b.</td>
</tr>
<tr>
<td>a &lt; b</td>
<td>boolean</td>
<td>The cube a is less than the cube b.</td>
</tr>
</tbody>
</table>
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>Operator</th>
<th>Result</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>a &lt;= b</td>
<td>boolean</td>
<td>The cube a is less than or equal to the cube b.</td>
</tr>
<tr>
<td>a &gt; b</td>
<td>boolean</td>
<td>The cube a is greater than the cube b.</td>
</tr>
<tr>
<td>a &gt;= b</td>
<td>boolean</td>
<td>The cube a is greater than or equal to the cube b.</td>
</tr>
<tr>
<td>a &lt;&gt; b</td>
<td>boolean</td>
<td>The cube a is not equal to the cube b.</td>
</tr>
<tr>
<td>a -&gt; n</td>
<td>float8</td>
<td>Get n-th coordinate of cube (counting from 1).</td>
</tr>
<tr>
<td>a ~&gt; n</td>
<td>float8</td>
<td>Get n-th coordinate of cube in following way: ( n = 2 \times k - 1 ) means lower bound of k-th dimension, ( n = 2 \times k ) means upper bound of k-th dimension. Negative n denotes the inverse value of the corresponding positive coordinate. This operator is designed for KNN-GiST support.</td>
</tr>
<tr>
<td>a &lt;-&gt; b</td>
<td>float8</td>
<td>Euclidean distance between a and b.</td>
</tr>
<tr>
<td>a &lt;#&gt; b</td>
<td>float8</td>
<td>Taxicab (L-1 metric) distance between a and b.</td>
</tr>
<tr>
<td>a &lt;=&gt; b</td>
<td>float8</td>
<td>Chebyshev (L-inf metric) distance between a and b.</td>
</tr>
</tbody>
</table>

(Before PostgreSQL 8.2, the containment operators @> and <@ were respectively called @ and ~. These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The scalar ordering operators (<, >=, etc) do not make a lot of sense for any practical purpose but sorting. These operators first compare the first coordinates, and if those are equal, compare the second coordinates, etc. They exist mainly to support the b-tree index operator class for cube, which can be useful for example if you would like a UNIQUE constraint on a cube column.

The cube module also provides a GiST index operator class for cube values. A cube GiST index can be used to search for values using the =, &, @>, and <@ operators in WHERE clauses.

In addition, a cube GiST index can be used to find nearest neighbors using the metric operators <->, <#>, and <=> in ORDER BY clauses. For example, the nearest neighbor of the 3-D point (0.5, 0.5, 0.5) could be found efficiently with:

```sql
SELECT c FROM test ORDER BY c <-> cube(array[0.5,0.5,0.5]) LIMIT 1;
```

The -> operator can also be used in this way to efficiently retrieve the first few values sorted by a selected coordinate. For example, to get the first few cubes ordered by the first coordinate (lower left corner) ascending one could use the following query:

```sql
SELECT c FROM test ORDER BY c -> 1 LIMIT 5;
```

And to get 2-D cubes ordered by the first coordinate of the upper right corner descending:
Table F.4. Cube Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Result</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>cube(float8)</code></td>
<td><code>cube</code></td>
<td>Makes a one dimensional cube with both coordinates the same.</td>
<td><code>cube(1) == '(1)'</code></td>
</tr>
<tr>
<td><code>cube(float8, float8)</code></td>
<td><code>cube</code></td>
<td>Makes a one dimensional cube.</td>
<td><code>cube(1,2) == '(1),(2)'</code></td>
</tr>
<tr>
<td><code>cube(float8[])</code></td>
<td><code>cube</code></td>
<td>Makes a zero-volume cube using the coordinates defined by the array.</td>
<td><code>cube(ARRAY[1,2]) == '(1,2)'</code></td>
</tr>
<tr>
<td><code>cube(float8[], float8[])</code></td>
<td><code>cube</code></td>
<td>Makes a cube with upper right and lower left coordinates as defined by the two arrays, which must be of the same length.</td>
<td><code>cube(ARRAY[1,2], ARRAY[3,4]) == '(1,2),(3,4)'</code></td>
</tr>
<tr>
<td><code>cube(cube, float8)</code></td>
<td><code>cube</code></td>
<td>Makes a new cube by adding a dimension on to an existing cube, with the same values for both endpoints of the new coordinate. This is useful for building cubes piece by piece from calculated values.</td>
<td><code>cube('(1,2), (3,4)', 5) == '(1,2,5), (3,4,5)'</code></td>
</tr>
<tr>
<td><code>cube(cube, float8, float8)</code></td>
<td><code>cube</code></td>
<td>Makes a new cube by adding a dimension on to an existing cube. This is useful for building cubes piece by piece from calculated values.</td>
<td><code>cube(cube('(1,2), (3,4)'), 5, 6) == '(1,2,5), (3,4,6)'</code></td>
</tr>
<tr>
<td><code>cube_dim(cube)</code></td>
<td><code>integer</code></td>
<td>Returns the number of dimensions of the cube.</td>
<td><code>cube_dim('(1,2), (3,4)') == 2</code></td>
</tr>
<tr>
<td><code>cube_ll_coord(cube, integer)</code></td>
<td><code>float8</code></td>
<td>Returns the n-th coordinate value for the lower left corner of the cube.</td>
<td><code>cube_ll_coord('(1,2), (3,4)'), 2) == 2'</code></td>
</tr>
<tr>
<td><code>cube_ur_coord(cube, integer)</code></td>
<td><code>float8</code></td>
<td>Returns the n-th coordinate value for the upper right corner of the cube.</td>
<td><code>cube_ur_coord('(1,2), (3,4)'), 2) == 4'</code></td>
</tr>
<tr>
<td><code>cube_is_point(cube)</code></td>
<td><code>boolean</code></td>
<td>Returns true if the cube is a point, that is, the two defining corners are the same.</td>
<td></td>
</tr>
<tr>
<td><code>cube_distance(cube, cube)</code></td>
<td><code>float8</code></td>
<td>Returns the distance between two cubes. If both cubes are points, this is the normal distance function.</td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Result</td>
<td>Description</td>
<td>Example</td>
</tr>
<tr>
<td>---------------------------</td>
<td>--------</td>
<td>-----------------------------------------------------------------------------------------------</td>
<td>----------------------------------------------</td>
</tr>
</tbody>
</table>
| cube_subset(cube, integer[]) | cube   | Makes a new cube from an existing cube, using a list of dimension indexes from an array. Can be used to extract the endpoints of a single dimension, or to drop dimensions, or to reorder them as desired. | cube_subset(cube('(1,3,5),(6,7,8)'), ARRAY[2]) == '(3),(7)'  
cube_subset(cube('(1,3,5),(6,7,8)'), ARRAY[3,2,1,1]) == '(5,3,1,1), (8,7,6,6)' |
| cube_union(cube, cube)     | cube   | Produces the union of two cubes.                                                              |                                              |
| cube_inter(cube, cube)     | cube   | Produces the intersection of two cubes.                                                       |                                              |
| cube_enlarge(cube, r double, n integer) | cube   | Increases the size of the cube by the specified radius \( r \) in at least \( n \) dimensions. If the radius is negative the cube is shrunk instead. All defined dimensions are changed by the radius \( r \). Lower-left coordinates are decreased by \( r \) and upper-right coordinates are increased by \( r \). If a lower-left coordinate is increased to more than the corresponding upper-right coordinate (this can only happen when \( r < 0 \)) than both coordinates are set to their average. If \( n \) is greater than the number of defined dimensions and the cube is being enlarged (\( r > 0 \)), then extra dimensions are added to make \( n \) altogether; 0 is used as the initial value for the extra coordinates. This function is useful for creating bounding boxes around a point for searching for nearby points. | cube_enlarge('(1,2),(3,4)', 0.5, 3) == '(0.5,1.5,-0.5), (3.5,4.5,0.5)' |

**F.9.4. Defaults**

I believe this union:

```sql
select cube_union('(0,5,2),(2,3,1)', '0');
cube_union
```
does not contradict common sense, neither does the intersection

\[
\text{select cube_inter('}(0,-1),(1,1)', '(-2),(2)');
\]

In all binary operations on differently-dimensioned cubes, I assume the lower-dimensional one to be a Cartesian projection, i.e., having zeroes in place of coordinates omitted in the string representation. The above examples are equivalent to:

\[
\text{cube_union('}(0,5,2),(2,3,1)', '}(0,0,0),(0,0,0)');
\]

The following containment predicate uses the point syntax, while in fact the second argument is internally represented by a box. This syntax makes it unnecessary to define a separate point type and functions for (box,point) predicates.

\[
\text{select cube_contains('}(0,0),(1,1)'', '0.5,0.5');
\]

F.9.5. Notes

For examples of usage, see the regression test sql/cube.sql.

To make it harder for people to break things, there is a limit of 100 on the number of dimensions of cubes. This is set in cubedata.h if you need something bigger.

F.9.6. Credits

Original author: Gene Selkov, Jr. <selkovjr@mcs.anl.gov>, Mathematics and Computer Science Division, Argonne National Laboratory.

My thanks are primarily to Prof. Joe Hellerstein (http://db.cs.berkeley.edu/jmh/) for elucidating the gist of the GiST (http://gist.cs.berkeley.edu/), and to his former student Andy Dong for his example written for Illustra. I am also grateful to all Postgres developers, present and past, for enabling myself to create my own world and live undisturbed in it. And I would like to acknowledge my gratitude to Argonne Lab and to the U.S. Department of Energy for the years of faithful support of my database research.

Minor updates to this package were made by Bruno Wolff III <bruno@wolff.to> in August/September of 2002. These include changing the precision from single precision to double precision and adding some new functions.

Additional updates were made by Joshua Reich <josh@root.net> in July 2006. These include \texttt{cube(float8[], float8[])} and cleaning up the code to use the V1 call protocol instead of the deprecated V0 protocol.
F.10. dblink

dblink is a module that supports connections to other PostgreSQL databases from within a database session.

See also postgres_fdw, which provides roughly the same functionality using a more modern and standards-compliant infrastructure.
**dblink_connect**

*dblink_connect* — opens a persistent connection to a remote database

**Synopsis**

```
dblink_connect(text connstr) returns text
dblink_connect(text connname, text connstr) returns text
```

**Description**

`dblink_connect()` establishes a connection to a remote PostgreSQL database. The server and database to be contacted are identified through a standard libpq connection string. Optionally, a name can be assigned to the connection. Multiple named connections can be open at once, but only one unnamed connection is permitted at a time. The connection will persist until closed or until the database session is ended.

The connection string may also be the name of an existing foreign server. It is recommended to use the foreign-data wrapper `dblink_fdw` when defining the foreign server. See the example below, as well as `CREATE SERVER` and `CREATE USER MAPPING`.

**Arguments**

*connname*

The name to use for this connection; if omitted, an unnamed connection is opened, replacing any existing unnamed connection.

*connstr*

libpq-style connection info string, for example `hostaddr=127.0.0.1 port=5432 dbname=mydb user=postgres password=mypassword options=-csearch_path=`. For details see Section 34.1.1. Alternatively, the name of a foreign server.

**Return Value**

Returns status, which is always `OK` (since any error causes the function to throw an error instead of returning).

**Notes**

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from `search_path`. One could, for example, add `options=-csearch_path=` to `connstr`. This consideration is not specific to dblink; it applies to every interface for executing arbitrary SQL commands.

Only superusers may use `dblink_connect` to create non-password-authenticated connections. If non-superusers need this capability, use `dblink_connect_u` instead.

It is unwise to choose connection names that contain equal signs, as this opens a risk of confusion with connection info strings in other dblink functions.

**Examples**

```
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
```

---

3314
SELECT dblink_connect('myconn', 'dbname=postgres options=-csearch_path=');
    dblink_connect
----------------
OK
(1 row)

-- FOREIGN DATA WRAPPER functionality
-- Note: local connection must require password authentication for
-- this to work properly
-- Otherwise, you will receive the following error from
dblink_connect():
--
-- ERROR: password is required
-- DETAIL: Non-superuser cannot connect if the server does
-- not request a password.
-- HINT: Target server's authentication method must be
-- changed.

CREATE SERVER fdtest FOREIGN DATA WRAPPER dblink_fdw OPTIONS
  (hostaddr '127.0.0.1', dbname 'contrib_regression');

CREATE USER regress_dblink_user WITH PASSWORD 'secret';
CREATE USER MAPPING FOR regress_dblink_user SERVER fdtest OPTIONS
  (user 'regress_dblink_user', password 'secret');
GRANT USAGE ON FOREIGN SERVER fdtest TO regress_dblink_user;
GRANT SELECT ON TABLE foo TO regress_dblink_user;

\set ORIGINAL_USER :USER
\c - regress_dblink_user
SELECT dblink_connect('myconn', 'fdtest');
    dblink_connect
----------------
OK
(1 row)

SELECT * FROM dblink('myconn','SELECT * FROM foo') AS t(a int, b
text, c text[]=);
a | b | c
------------------------
0  | a | {a0,b0,c0}
1  | b | {a1,b1,c1}
2  | c | {a2,b2,c2}
3  | d | {a3,b3,c3}
4  | e | {a4,b4,c4}
5  | f | {a5,b5,c5}
6  | g | {a6,b6,c6}
7  | h | {a7,b7,c7}
8  | i | {a8,b8,c8}
9  | j | {a9,b9,c9}
10 | k | {a10,b10,c10}
(11 rows)
\c - :ORIGINAL_USER
REVOKE USAGE ON FOREIGN SERVER fdtest FROM regress_dblink_user;
REVOKE SELECT ON TABLE foo FROM regress_dblink_user;
DROP USER MAPPING FOR regress_dblink_user SERVER fdtest;
DROP USER regress_dblink_user;
DROP SERVER fdtest;
**dblink_connect_u**

`dblink_connect_u` — opens a persistent connection to a remote database, insecurely

**Synopsis**

```plaintext
dblink_connect_u(text connstr) returns text
dblink_connect_u(text connname, text connstr) returns text
```

**Description**

`dblink_connect_u()` is identical to `dblink_connect()`, except that it will allow non-superusers to connect using any authentication method.

If the remote server selects an authentication method that does not involve a password, then impersonation and subsequent escalation of privileges can occur, because the session will appear to have originated from the user as which the local PostgreSQL server runs. Also, even if the remote server does demand a password, it is possible for the password to be supplied from the server environment, such as a `~/.pgpass` file belonging to the server's user. This opens not only a risk of impersonation, but the possibility of exposing a password to an untrustworthy remote server. Therefore, `dblink_connect_u()` is initially installed with all privileges revoked from `PUBLIC`, making it un-callable except by superusers. In some situations it may be appropriate to grant `EXECUTE` permission for `dblink_connect_u()` to specific users who are considered trustworthy, but this should be done with care. It is also recommended that any `~/.pgpass` file belonging to the server's user not contain any records specifying a wildcard host name.

For further details see `dblink_connect()`.
**dblink_disconnect**

dblink_disconnect — closes a persistent connection to a remote database

**Synopsis**

dblink_disconnect() returns text
dblink_disconnect(text conname) returns text

**Description**

dblink_disconnect() closes a connection previously opened by dblink_connect(). The form with no arguments closes an unnamed connection.

**Arguments**

conname

The name of a named connection to be closed.

**Return Value**

Returns status, which is always OK (since any error causes the function to throw an error instead of returning).

**Examples**

```sql
SELECT dblink_disconnect();
dblink_disconnect
-------------------
OK
(1 row)

SELECT dblink_disconnect('myconn');
dblink_disconnect
-------------------
OK
(1 row)
```
**dblink**

dblink — executes a query in a remote database

**Synopsis**

\[
\begin{align*}
\text{dblink(text connname, text sql [, bool fail_on_error]) returns setof record} \\
\text{dlink(text connstr, text sql [, bool fail_on_error]) returns setof record} \\
\text{dlink(text sql [, bool fail_on_error]) returns setof record}
\end{align*}
\]

**Description**

dblink executes a query (usually a SELECT, but it can be any SQL statement that returns rows) in a remote database.

When two text arguments are given, the first one is first looked up as a persistent connection's name; if found, the command is executed on that connection. If not found, the first argument is treated as a connection info string as for dblink_connect, and the indicated connection is made just for the duration of this command.

**Arguments**

- **connname**
  
  Name of the connection to use; omit this parameter to use the unnamed connection.

- **connstr**
  
  A connection info string, as previously described for dblink_connect.

- **sql**
  
  The SQL query that you wish to execute in the remote database, for example `select * from foo`.

- **fail_on_error**
  
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

The function returns the row(s) produced by the query. Since dblink can be used with any query, it is declared to return record, rather than specifying any particular set of columns. This means that you must specify the expected set of columns in the calling query — otherwise PostgreSQL would not know what to expect. Here is an example:

```sql
SELECT *
  FROM dblink('dbname=mydb options=-csearch_path=',
               'select proname, prosrc from pg_proc')
               AS t1(proname name, prosrc text)
WHERE proname LIKE 'bytea%';
```
The “alias” part of the FROM clause must specify the column names and types that the function will return. (Specifying column names in an alias is actually standard SQL syntax, but specifying column types is a PostgreSQL extension.) This allows the system to understand what * should expand to, and what proname in the WHERE clause refers to, in advance of trying to execute the function. At runtime, an error will be thrown if the actual query result from the remote database does not have the same number of columns shown in the FROM clause. The column names need not match, however, and dblink does not insist on exact type matches either. It will succeed so long as the returned data strings are valid input for the column type declared in the FROM clause.

Notes

A convenient way to use dblink with predetermined queries is to create a view. This allows the column type information to be buried in the view, instead of having to spell it out in every query. For example,

```
CREATE VIEW myremote_pg_proc AS
  SELECT *
    FROM dblink('dbname=postgres options=-csearch_path=','select proname, prosrc from pg_proc')
    AS t1(proname name, prosrc text);
```

```
SELECT * FROM myremote_pg_proc WHERE proname LIKE 'bytea%';
```

Examples

```
SELECT * FROM dblink('dbname=postgres options=-csearch_path=','select proname, prosrc from pg_proc')
    AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';
```

<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>

(12 rows)

```
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
```

dblink_connect

<table>
<thead>
<tr>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>OK</td>
</tr>
</tbody>
</table>

(1 row)

```
SELECT * FROM dblink('select proname, prosrc from pg_proc')
    AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';
```

<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
</tbody>
</table>

3320
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>module</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytealt</td>
<td></td>
</tr>
<tr>
<td>byteale</td>
<td></td>
</tr>
<tr>
<td>byteagt</td>
<td></td>
</tr>
<tr>
<td>byteage</td>
<td></td>
</tr>
<tr>
<td>byteane</td>
<td></td>
</tr>
<tr>
<td>byteacmp</td>
<td></td>
</tr>
<tr>
<td>bytealike</td>
<td></td>
</tr>
<tr>
<td>byteanlike</td>
<td></td>
</tr>
<tr>
<td>byteain</td>
<td></td>
</tr>
<tr>
<td>byteaout</td>
<td></td>
</tr>
</tbody>
</table>

(12 rows)

```sql
SELECT dblink_connect('myconn', 'dbname=regression options=- csearch_path=''');
```

```sql
OK
```

(1 row)

```sql
SELECT * FROM dblink('myconn', 'select proname, prosrc from pg_proc')
  AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';
```

<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytearecv</td>
<td>bytearecv</td>
</tr>
<tr>
<td>byteasend</td>
<td>byteasend</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>

(14 rows)
**dblink_exec**

*dblink_exec* — executes a command in a remote database

**Synopsis**

```plaintext
dblink_exec(text connname, text sql [, bool fail_on_error]) returns text
dblink_exec(text connstr, text sql [, bool fail_on_error]) returns text
dblink_exec(text sql [, bool fail_on_error]) returns text
```

**Description**

`dblink_exec` executes a command (that is, any SQL statement that doesn't return rows) in a remote database.

When two `text` arguments are given, the first one is first looked up as a persistent connection's name; if found, the command is executed on that connection. If not found, the first argument is treated as a connection info string as for `dblink_connect`, and the indicated connection is made just for the duration of this command.

**Arguments**

- **connname**
  
  Name of the connection to use; omit this parameter to use the unnamed connection.

- **connstr**
  
  A connection info string, as previously described for `dblink_connect`.

- **sql**
  
  The SQL command that you wish to execute in the remote database, for example `insert into foo values(0,'a',"{"a0","b0","c0"}")`.

- **fail_on_error**
  
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function's return value is set to `ERROR`.

**Return Value**

Returns status, either the command's status string or `ERROR`.

**Examples**

```sql
SELECT dblink_connect('dbname=dblink_test_standby');
  dblink_connect
------------------
  OK
(1 row)
```
SELECT dblink_exec('insert into foo
values(21,''z'',''{"a0","b0","c0"}"');');
    dblink_exec
-----------------
    INSERT 943366 1
    (1 row)

SELECT dblink_connect('myconn', 'dbname=regression');
    dblink_connect
-------------
    OK
    (1 row)

SELECT dblink_exec('myconn', 'insert into foo
values(21,''z'',''{"a0","b0","c0"}"');');
    dblink_exec
------------------
    INSERT 6432584 1
    (1 row)

SELECT dblink_exec('myconn', 'insert into pg_class values
(''foo'')',false);
NOTICE: sql error
DETAIL: ERROR: null value in column "relnamespace" violates not-null constraint
    dblink_exec
-------------
    ERROR
    (1 row)
**dblink_open**

`dblink_open` — opens a cursor in a remote database

### Synopsis

```c

$dblink_open$(text $cursorname$, text $sql$ [, bool $fail_on_error$])
returns text
$dblink_open$(text $connname$, text $cursorname$, text $sql$ [, bool $fail_on_error$])
returns text
```

### Description

`dblink_open()` opens a cursor in a remote database. The cursor can subsequently be manipulated with `dblink_fetch()` and `dblink_close()`.

### Arguments

- **$connname$**
  - Name of the connection to use; omit this parameter to use the unnamed connection.

- **$cursorname$**
  - The name to assign to this cursor.

- **$sql$**
  - The `SELECT` statement that you wish to execute in the remote database, for example `select * from pg_class`.

- **$fail_on_error$**
  - If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function's return value is set to `ERROR`.

### Return Value

Returns status, either `OK` or `ERROR`.

### Notes

Since a cursor can only persist within a transaction, `dblink_open` starts an explicit transaction block (`BEGIN`) on the remote side, if the remote side was not already within a transaction. This transaction will be closed again when the matching `dblink_close` is executed. Note that if you use `dblink_exec` to change data between `dblink_open` and `dblink_close`, and then an error occurs or you use `dblink_disconnect` before `dblink_close`, your change will be lost because the transaction will be aborted.

### Examples

```c
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
dblink_connect
```
---
OK
(1 row)

SELECT dblink_open('foo', 'select proname, prosrc from pg_proc');
dblink_open
---
OK
(1 row)
**dblink_fetch**

dblink_fetch — returns rows from an open cursor in a remote database

**Synopsis**

```
dblink_fetch(text cursorname, int howmany [, bool fail_on_error]) returns setof record
```

```
dblink_fetch(text connname, text cursorname, int howmany [, bool fail_on_error]) returns setof record
```

**Description**

`dblink_fetch` fetches rows from a cursor previously established by `dblink_open`.

**Arguments**

- **connname**
  Name of the connection to use; omit this parameter to use the unnamed connection.

- **cursorname**
  The name of the cursor to fetch from.

- **howmany**
  The maximum number of rows to retrieve. The next `howmany` rows are fetched, starting at the current cursor position, moving forward. Once the cursor has reached its end, no more rows are produced.

- **fail_on_error**
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

The function returns the row(s) fetched from the cursor. To use this function, you will need to specify the expected set of columns, as previously discussed for `dblink`.

**Notes**

On a mismatch between the number of return columns specified in the FROM clause, and the actual number of columns returned by the remote cursor, an error will be thrown. In this event, the remote cursor is still advanced by as many rows as it would have been if the error had not occurred. The same is true for any other error occurring in the local query after the remote FETCH has been done.

**Examples**

```sql
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
dblink_connect
-------------
```
SELECT dblink_open('foo', 'select proname, prosrc from pg_proc where proname like ''bytea%''');

-------------
OK
(1 row)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);

<table>
<thead>
<tr>
<th>funcname</th>
<th>source</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
</tbody>
</table>

(5 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);

<table>
<thead>
<tr>
<th>funcname</th>
<th>source</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
</tbody>
</table>

(5 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);

<table>
<thead>
<tr>
<th>funcname</th>
<th>source</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>

(2 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);

<table>
<thead>
<tr>
<th>funcname</th>
<th>source</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>

(0 rows)
**dblink_close**

`dblink_close` — closes a cursor in a remote database

### Synopsis

```
dblink_close(text cursorname [, bool fail_on_error]) returns text
```

```
dblink_close(text connname, text cursorname [, bool fail_on_error]) returns text
```

### Description

`dblink_close` closes a cursor previously opened with `dblink_open`.

### Arguments

- `connname`
  
  Name of the connection to use; omit this parameter to use the unnamed connection.

- `cursorname`
  
  The name of the cursor to close.

- `fail_on_error`
  
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function's return value is set to `ERROR`.

### Return Value

Returns status, either `OK` or `ERROR`.

### Notes

If `dblink_open` started an explicit transaction block, and this is the last remaining open cursor in this connection, `dblink_close` will issue the matching `COMMIT`.

### Examples

```
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
   dblink_connect
----------------
    OK
(1 row)

SELECT dblink_open('foo', 'select proname, prosrc from pg_proc');
   dblink_open
----------
    OK
(1 row)

SELECT dblink_close('foo');
```
Additional Supplied Modules

dblink_close
------------
OK
(1 row)
**dblink_get_connections**

*dblink_get_connections* — returns the names of all open named dblink connections

**Synopsis**

```
dblink_get_connections() returns text[]
```

**Description**

`dblink_get_connections` returns an array of the names of all open named dblink connections.

**Return Value**

Returns a text array of connection names, or NULL if none.

**Examples**

```
SELECT dblink_get_connections();
```
**dblink_error_message**

`dblink_error_message` — gets last error message on the named connection

**Synopsis**

```
`dblink_error_message(text conname) returns text`
```

**Description**

*dblink_error_message* fetches the most recent remote error message for a given connection.

**Arguments**

`conname`

Name of the connection to use.

**Return Value**

Returns last error message, or an empty string if there has been no error in this connection.

**Examples**

```
SELECT dblink_error_message('dtest1');
```
**dblink_send_query**

*dblink_send_query* — sends an async query to a remote database

**Synopsis**

```c
dblink_send_query(text connname, text sql) returns int
```

**Description**

`dblink_send_query` sends a query to be executed asynchronously, that is, without immediately waiting for the result. There must not be an async query already in progress on the connection.

After successfully dispatching an async query, completion status can be checked with `dblink_is_busy`, and the results are ultimately collected with `dblink_get_result`. It is also possible to attempt to cancel an active async query using `dblink_cancel_query`.

**Arguments**

* `connname`
  
  Name of the connection to use.

* `sql`
  
  The SQL statement that you wish to execute in the remote database, for example `select * from pg_class`.

**Return Value**

Returns 1 if the query was successfully dispatched, 0 otherwise.

**Examples**

```sql
SELECT dblink_send_query('dtest1', 'SELECT * FROM foo WHERE f1 < 3');
```
**dblink_is_busy**

`dblink_is_busy` — checks if connection is busy with an async query

**Synopsis**

```plaintext
dblink_is_busy(text conname) returns int
```

**Description**

`dblink_is_busy` tests whether an async query is in progress.

**Arguments**

`conname`

Name of the connection to check.

**Return Value**

Returns 1 if connection is busy, 0 if it is not busy. If this function returns 0, it is guaranteed that `dblink_get_result` will not block.

**Examples**

```sql
SELECT dblink_is_busy('dtest1');
```
dblink_get_notify

dblink_get_notify — retrieve async notifications on a connection

Synopsis

dblink_get_notify() returns setof (notify_name text, be_pid int, extra text)
dblink_get_notify(text connname) returns setof (notify_name text, be_pid int, extra text)

Description

dblink_get_notify retrieves notifications on either the unnamed connection, or on a named connection if specified. To receive notifications via dblink, LISTEN must first be issued, using dblink_exec. For details see LISTEN and NOTIFY.

Arguments

connname

The name of a named connection to get notifications on.

Return Value

Returns setof (notify_name text, be_pid int, extra text), or an empty set if none.

Examples

SELECT dblink_exec('LISTEN virtual');
dblink_exec

-------------
LISTEN
(1 row)

SELECT * FROM dblink_get_notify();

notify_name | be_pid | extra
-----------------------------
(0 rows)

NOTIFY virtual;
NOTIFY

SELECT * FROM dblink_get_notify();

notify_name | be_pid | extra
-----------------------------
virtual | 1229 |
(1 row)
Additional Supplied Modules

**dblink_get_result**

`dblink_get_result` — gets an async query result

**Synopsis**

```
dblink_get_result(text connname [, bool fail_on_error]) returns setof record
```

**Description**

`dblink_get_result` collects the results of an asynchronous query previously sent with `dblink_send_query`. If the query is not already completed, `dblink_get_result` will wait until it is.

**Arguments**

- `connname`
  - Name of the connection to use.
- `fail_on_error`
  - If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

For an async query (that is, a SQL statement returning rows), the function returns the row(s) produced by the query. To use this function, you will need to specify the expected set of columns, as previously discussed for `dblink`.

For an async command (that is, a SQL statement not returning rows), the function returns a single row with a single text column containing the command's status string. It is still necessary to specify that the result will have a single text column in the calling `FROM` clause.

**Notes**

This function must be called if `dblink_send_query` returned 1. It must be called once for each query sent, and one additional time to obtain an empty set result, before the connection can be used again.

When using `dblink_send_query` and `dblink_get_result`, dblink fetches the entire remote query result before returning any of it to the local query processor. If the query returns a large number of rows, this can result in transient memory bloat in the local session. It may be better to open such a query as a cursor with `dblink_open` and then fetch a manageable number of rows at a time. Alternatively, use plain `dblink()`, which avoids memory bloat by spooling large result sets to disk.

**Examples**

```contrib_regression=# SELECT dblink_connect('dtest1', 'dbname=contrib_regression');
dblink_connect
----------------
3335```
Additional Supplied Modules

```
OK
(1 row)

contrib_regression=# SELECT * FROM
dblink_send_query('dtest1', 'select * from foo
where f1 < 3') AS t1;
t1----
1
(1 row)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS
t1(f1 int, f2 text, f3 text[]);
f1 | f2 |     f3
----+----+------------
 0 | a  | {a0,b0,c0}
 1 | b  | {a1,b1,c1}
 2 | c  | {a2,b2,c2}
(3 rows)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS
t1(f1 int, f2 text, f3 text[]);
f1 | f2 |     f3
----+----+------------
 0 | a  | {a0,b0,c0}
 1 | b  | {a1,b1,c1}
 2 | c  | {a2,b2,c2}
(3 rows)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS
t1(f1 int, f2 text, f3 text[]);
f1 | f2 |     f3
----+----+------------
 0 | a  | {a0,b0,c0}
 1 | b  | {a1,b1,c1}
 2 | c  | {a2,b2,c2}
(3 rows)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS
t1(f1 int, f2 text, f3 text[]);
f1 | f2 |     f3
----+----+------------
 7 | h  | {a7,b7,c7}
 8 | i  | {a8,b8,c8}
 9 | j  | {a9,b9,c9}
10 | k  | {a10,b10,c10}
(4 rows)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS
t1(f1 int, f2 text, f3 text[]);
f1 | f2 |     f3
----+----+------------
(0 rows)
```

3336
**dblink_cancel_query**

*dblink_cancel_query* — cancels any active query on the named connection

### Synopsis

```
dblink_cancel_query(text connname) returns text
```

### Description

`dblink_cancel_query` attempts to cancel any query that is in progress on the named connection. Note that this is not certain to succeed (since, for example, the remote query might already have finished). A cancel request simply improves the odds that the query will fail soon. You must still complete the normal query protocol, for example by calling `dblink_get_result`.

### Arguments

* `connname`
  
  Name of the connection to use.

### Return Value

Returns `OK` if the cancel request has been sent, or the text of an error message on failure.

### Examples

```
SELECT dblink_cancel_query('dtest1');
```
**dblink_get_pkey**

`dblink_get_pkey` — returns the positions and field names of a relation's primary key fields

**Synopsis**

`dblink_get_pkey(text relname) returns setof dblink_pkey_results`

**Description**

`dblink_get_pkey` provides information about the primary key of a relation in the local database. This is sometimes useful in generating queries to be sent to remote databases.

**Arguments**

`relname`

Name of a local relation, for example `foo` or `myschema.mytab`. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

**Return Value**

Returns one row for each primary key field, or no rows if the relation has no primary key. The result row type is defined as

```sql
CREATE TYPE dblink_pkey_results AS (position int, colname text);
```

The `position` column simply runs from 1 to \( N \); it is the number of the field within the primary key, not the number within the table's columns.

**Examples**

```sql
CREATE TABLE foobar (  
    f1 int,  
    f2 int,  
    f3 int,  
    PRIMARY KEY (f1, f2, f3)  
);  
CREATE TABLE

SELECT * FROM dblink_get_pkey('foobar');

<table>
<thead>
<tr>
<th>position</th>
<th>colname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>f1</td>
</tr>
<tr>
<td>2</td>
<td>f2</td>
</tr>
<tr>
<td>3</td>
<td>f3</td>
</tr>
</tbody>
</table>
```

(3 rows)
Additional Supplied Modules

**dblink_build_sql_insert**

`dblink_build_sql_insert` — builds an INSERT statement using a local tuple, replacing the primary key field values with alternative supplied values

**Synopsis**

```sql
dblink_build_sql_insert(text relname,
                        int2vector primary_key_attnums,
                        integer num_primary_key_atts,
                        text[] src_pk_att_vals_array,
                        text[] tgt_pk_att_vals_array) returns text
```

**Description**

`dblink_build_sql_insert` can be useful in doing selective replication of a local table to a remote database. It selects a row from the local table based on primary key, and then builds a SQL INSERT command that will duplicate that row, but with the primary key values replaced by the values in the last argument. (To make an exact copy of the row, just specify the same values for the last two arguments.)

**Arguments**

- `relname`
  
  Name of a local relation, for example `foo` or `myschema.mytab`. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

- `primary_key_attnums`
  
  Attribute numbers (1-based) of the primary key fields, for example `1 2`.

- `num_primary_key_atts`
  
  The number of primary key fields.

- `src_pk_att_vals_array`
  
  Values of the primary key fields to be used to look up the local tuple. Each field is represented in text form. An error is thrown if there is no local row with these primary key values.

- `tgt_pk_att_vals_array`
  
  Values of the primary key fields to be placed in the resulting INSERT command. Each field is represented in text form.

**Return Value**

Returns the requested SQL statement as text.

**Notes**

As of PostgreSQL 9.0, the attribute numbers in `primary_key_attnums` are interpreted as logical column numbers, corresponding to the column's position in `SELECT * FROM relname`. Previous
versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

Examples

```
SELECT dblink_build_sql_insert('foo', '1 2', 2, '{"1", "a"}', '{"1", "b''a"}');

-----------------------------
dblink_build_sql_insert

INSERT INTO foo(f1,f2,f3) VALUES('1','b''a','1')
(1 row)
```
**dblink_build_sql_delete**

`dblink_build_sql_delete` — builds a DELETE statement using supplied values for primary key fields.

### Synopsis

```
dblink_build_sql_delete(text relname, int2vector primary_key_attnums, integer num_primary_key_atts, text[] tgt_pk_att_vals_array) returns text
```

### Description

`dblink_build_sql_delete` can be useful in doing selective replication of a local table to a remote database. It builds a SQL DELETE command that will delete the row with the given primary key values.

### Arguments

- **relname**
  Name of a local relation, for example `foo` or `myschema.mytab`. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

- **primary_key_attnums**
  Attribute numbers (1-based) of the primary key fields, for example `1 2`.

- **num_primary_key_atts**
  The number of primary key fields.

- **tgt_pk_att_vals_array**
  Values of the primary key fields to be used in the resulting DELETE command. Each field is represented in text form.

### Return Value

Returns the requested SQL statement as text.

### Notes

As of PostgreSQL 9.0, the attribute numbers in `primary_key_attnums` are interpreted as logical column numbers, corresponding to the column's position in `SELECT * FROM relname`. Previous versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

### Examples

```
SELECT dblink_build_sql_delete('"MyFoo"', '1 2', 2, '{"1", "b"}');
```

3341
DELETE FROM "MyFoo" WHERE f1='1' AND f2='b'
(1 row)
**dblink_build_sql_update**

dblink_build_sql_update — builds an UPDATE statement using a local tuple, replacing the primary key field values with alternative supplied values

**Synopsis**

dblink_build_sql_update(text relname,
int2vector primary_key_attnums,
integer num_primary_key_atts,
text[] src_pk_att_vals_array,
text[] tgt_pk_att_vals_array) returns text

**Description**

dblink_build_sql_update can be useful in doing selective replication of a local table to a remote database. It selects a row from the local table based on primary key, and then builds a SQL UPDATE command that will duplicate that row, but with the primary key values replaced by the values in the last argument. (To make an exact copy of the row, just specify the same values for the last two arguments.) The UPDATE command always assigns all fields of the row — the main difference between this and dblink_build_sql_insert is that it’s assumed that the target row already exists in the remote table.

**Arguments**

**rename**

Name of a local relation, for example foo or myschema.mytab. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

**primary_key_attnums**

Attribute numbers (1-based) of the primary key fields, for example 1 2.

**num_primary_key_atts**

The number of primary key fields.

**src_pk_att_vals_array**

Values of the primary key fields to be used to look up the local tuple. Each field is represented in text form. An error is thrown if there is no local row with these primary key values.

**tgt_pk_att_vals_array**

Values of the primary key fields to be placed in the resulting UPDATE command. Each field is represented in text form.

**Return Value**

Returns the requested SQL statement as text.

**Notes**

As of PostgreSQL 9.0, the attribute numbers in primary_key_attnums are interpreted as logical column numbers, corresponding to the column's position in SELECT * FROM rename. Previous
versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

Examples

```
SELECT dblink_build_sql_update('foo', '1 2', 2, '{"1", "a"}', '{"1", "b"}');
```

```
dblink_build_sql_update
------------------------------------------
UPDATE foo SET f1='1',f2='b',f3='1' WHERE f1='1' AND f2='b'
(1 row)
```

F.11. dict_int

dict_int is an example of an add-on dictionary template for full-text search. The motivation for this example dictionary is to control the indexing of integers (signed and unsigned), allowing such numbers to be indexed while preventing excessive growth in the number of unique words, which greatly affects the performance of searching.

F.11.1. Configuration

The dictionary accepts two options:

- The `maxlen` parameter specifies the maximum number of digits allowed in an integer word. The default value is 6.
- The `rejectlong` parameter specifies whether an overlength integer should be truncated or ignored. If `rejectlong` is `false` (the default), the dictionary returns the first `maxlen` digits of the integer. If `rejectlong` is `true`, the dictionary treats an overlength integer as a stop word, so that it will not be indexed. Note that this also means that such an integer cannot be searched for.

F.11.2. Usage

Installing the `dict_int` extension creates a text search template `intdict_template` and a dictionary `intdict` based on it, with the default parameters. You can alter the parameters, for example

```
mydb# ALTER TEXT SEARCH DICTIONARY intdict (MAXLEN = 4, REJECTLONG = true);
ALTER TEXT SEARCH DICTIONARY
```

or create new dictionaries based on the template.

To test the dictionary, you can try

```
mydb# select ts_lexize('intdict', '12345678');
```

```
ts_lexize
------
{123456}
```

but real-world usage will involve including it in a text search configuration as described in Chapter 12. That might look like this:

```
ALTER TEXT SEARCH CONFIGURATION english
    ALTER MAPPING FOR int, uint WITH intdict;
```
F.12. dict_xsyn

dict_xsyn (Extended Synonym Dictionary) is an example of an add-on dictionary template for full-text search. This dictionary type replaces words with groups of their synonyms, and so makes it possible to search for a word using any of its synonyms.

F.12.1. Configuration

A dict_xsyn dictionary accepts the following options:

- `matchorig` controls whether the original word is accepted by the dictionary. Default is `true`.
- `matchsynonyms` controls whether the synonyms are accepted by the dictionary. Default is `false`.
- `keeporig` controls whether the original word is included in the dictionary's output. Default is `true`.
- `keepsynonyms` controls whether the synonyms are included in the dictionary's output. Default is `true`.
- `rules` is the base name of the file containing the list of synonyms. This file must be stored in `$SHAREDIR/tsearch_data/` (where `$SHAREDIR` means the PostgreSQL installation's shared-data directory). Its name must end in `.rules` (which is not to be included in the `rules` parameter).

The rules file has the following format:

- Each line represents a group of synonyms for a single word, which is given first on the line. Synonyms are separated by whitespace, thus:

  ```
  word syn1 syn2 syn3
  ```

- The sharp (`#`) sign is a comment delimiter. It may appear at any position in a line. The rest of the line will be skipped.

Look at `xsyn_sample.rules`, which is installed in `$SHAREDIR/tsearch_data/`, for an example.

F.12.2. Usage

Installing the dict_xsyn extension creates a text search template `xsyn_template` and a dictionary `xsyn` based on it, with default parameters. You can alter the parameters, for example

```
mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules', KEEPORIG=false);
```

or create new dictionaries based on the template.

To test the dictionary, you can try

```
mydb=# SELECT ts_lexize('xsyn', 'word');
```
Additional Supplied Modules

{syn1,syn2,syn3}

mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules',
KEEPORIG=true);
ALTER TEXT SEARCH DICTIONARY

mydb=# SELECT ts_lexize('xsyn', 'word');
  ts_lexize
  -----------------------
  {word,syn1,syn2,syn3}

mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules',
KEEPORIG=false, MATCHSYNONYMS=true);
ALTER TEXT SEARCH DICTIONARY

mydb=# SELECT ts_lexize('xsyn', 'syn1');
  ts_lexize
  -----------------------
  {syn1,syn2,syn3}

mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules',
KEEPORIG=true, MATCHORIG=false, KEEPSYNONYMS=false);
ALTER TEXT SEARCH DICTIONARY

mydb=# SELECT ts_lexize('xsyn', 'syn1');
  ts_lexize
  -----------------------
  {word}

Real-world usage will involve including it in a text search configuration as described in Chapter 12. That might look like this:

ALTER TEXT SEARCH CONFIGURATION english
  ALTER MAPPING FOR word, asciiword WITH xsyn, english_stem;

F.13. earthdistance

The earthdistance module provides two different approaches to calculating great circle distances on the surface of the Earth. The one described first depends on the cube module (which must be installed before earthdistance can be installed). The second one is based on the built-in point data type, using longitude and latitude for the coordinates.

In this module, the Earth is assumed to be perfectly spherical. (If that's too inaccurate for you, you might want to look at the PostGIS\(^1\) project.)

F.13.1. Cube-based Earth Distances

Data is stored in cubes that are points (both corners are the same) using 3 coordinates representing the x, y, and z distance from the center of the Earth. A domain earth over cube is provided, which includes constraint checks that the value meets these restrictions and is reasonably close to the actual surface of the Earth.

The radius of the Earth is obtained from the earth() function. It is given in meters. But by changing this one function you can change the module to use some other units, or to use a different value of the radius that you feel is more appropriate.

\(^1\) http://postgis.net/
This package has applications to astronomical databases as well. Astronomers will probably want to change \texttt{earth()} to return a radius of $180/\pi()$ so that distances are in degrees.

Functions are provided to support input in latitude and longitude (in degrees), to support output of latitude and longitude, to calculate the great circle distance between two points and to easily specify a bounding box usable for index searches.

The provided functions are shown in Table F.5.

\textbf{Table F.5. Cube-based Earthdistance Functions}

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{earth()}</td>
<td>float8</td>
<td>Returns the assumed radius of the Earth.</td>
</tr>
<tr>
<td>\texttt{sec_to_gc(float8)}</td>
<td>float8</td>
<td>Converts the normal straight line (secant) distance between two points on the surface of the Earth to the great circle distance between them.</td>
</tr>
<tr>
<td>\texttt{gc_to_sec(float8)}</td>
<td>float8</td>
<td>Converts the great circle distance between two points on the surface of the Earth to the normal straight line (secant) distance between them.</td>
</tr>
<tr>
<td>\texttt{ll_to_earth(float8, float8)}</td>
<td>earth</td>
<td>Returns the location of a point on the surface of the Earth given its latitude (argument 1) and longitude (argument 2) in degrees.</td>
</tr>
<tr>
<td>\texttt{latitude(earth)}</td>
<td>float8</td>
<td>Returns the latitude in degrees of a point on the surface of the Earth.</td>
</tr>
<tr>
<td>\texttt{longitude(earth)}</td>
<td>float8</td>
<td>Returns the longitude in degrees of a point on the surface of the Earth.</td>
</tr>
<tr>
<td>\texttt{earth_distance(earth, earth)}</td>
<td>float8</td>
<td>Returns the great circle distance between two points on the surface of the Earth.</td>
</tr>
<tr>
<td>\texttt{earth_box(earth, float8)}</td>
<td>cube</td>
<td>Returns a box suitable for an indexed search using the cube @&gt; operator for points within a given great circle distance of a location. Some points in this box are further than the specified great circle distance from the location, so a second check using \texttt{earth_distance} should be included in the query.</td>
</tr>
</tbody>
</table>

\textbf{F.13.2. Point-based Earth Distances}

The second part of the module relies on representing Earth locations as values of type \texttt{point}, in which the first component is taken to represent longitude in degrees, and the second component is taken to represent latitude in degrees. Points are taken as (longitude, latitude) and not vice versa because longitude is closer to the intuitive idea of x-axis and latitude to y-axis.

A single operator is provided, shown in Table F.6.
### Table F.6. Point-based Earthdistance Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>point $&lt;$@$&gt;$ point</td>
<td>float8</td>
<td>Gives the distance in statute miles between two points on the Earth's surface.</td>
</tr>
</tbody>
</table>

Note that unlike the cube-based part of the module, units are hardwired here: changing the `earth()` function will not affect the results of this operator.

One disadvantage of the longitude/latitude representation is that you need to be careful about the edge conditions near the poles and near +/- 180 degrees of longitude. The cube-based representation avoids these discontinuities.

### F.14. file_fdw

The `file_fdw` module provides the foreign-data wrapper `file_fdw`, which can be used to access data files in the server's file system, or to execute programs on the server and read their output. The data file or program output must be in a format that can be read by `COPY FROM`; see COPY for details.

Access to data files is currently read-only.

A foreign table created using this wrapper can have the following options:

- **filename**
  - Specifies the file to be read. Must be an absolute path name. Either filename or program must be specified, but not both.

- **program**
  - Specifies the command to be executed. The standard output of this command will be read as though `COPY FROM PROGRAM` were used. Either program or filename must be specified, but not both.

- **format**
  - Specifies the data format, the same as COPY’s FORMAT option.

- **header**
  - Specifies whether the data has a header line, the same as COPY’s HEADER option.

- **delimiter**
  - Specifies the data delimiter character, the same as COPY’s DELIMITER option.

- **quote**
  - Specifies the data quote character, the same as COPY’s QUOTE option.

- **escape**
  - Specifies the data escape character, the same as COPY’s ESCAPE option.

- **null**
  - Specifies the data null string, the same as COPY’s NULL option.

- **encoding**
  - Specifies the data encoding, the same as COPY’s ENCODING option.
Note that while COPY allows options such as HEADER to be specified without a corresponding value, the foreign table option syntax requires a value to be present in all cases. To activate COPY options typically written without a value, you can pass the value TRUE, since all such options are Booleans.

A column of a foreign table created using this wrapper can have the following options:

**force_not_null**

This is a Boolean option. If true, it specifies that values of the column should not be matched against the null string (that is, the table-level null option). This has the same effect as listing the column in COPY's FORCE_NOT_NULL option.

**force_null**

This is a Boolean option. If true, it specifies that values of the column which match the null string are returned as NULL even if the value is quoted. Without this option, only unquoted values matching the null string are returned as NULL. This has the same effect as listing the column in COPY's FORCE_NULL option.

COPY's OIDS and FORCE_QUOTE options are currently not supported by file_fdw.

These options can only be specified for a foreign table or its columns, not in the options of the file_fdw foreign-data wrapper, nor in the options of a server or user mapping using the wrapper.

Changing table-level options requires being a superuser or having the privileges of the default role pg_read_server_files (to use a filename) or the default role pg_execute_server_programs (to use a program), for security reasons: only certain users should be able to control which file is read or which program is run. In principle regular users could be allowed to change the other options, but that's not supported at present.

When specifying the program option, keep in mind that the option string is executed by the shell. If you need to pass any arguments to the command that come from an untrusted source, you must be careful to strip or escape any characters that might have special meaning to the shell. For security reasons, it is best to use a fixed command string, or at least avoid passing any user input in it.

For a foreign table using file_fdw, EXPLAIN shows the name of the file to be read or program to be run. For a file, unless COSTS OFF is specified, the file size (in bytes) is shown as well.

**Example F.1. Create a Foreign Table for PostgreSQL CSV Logs**

One of the obvious uses for file_fdw is to make the PostgreSQL activity log available as a table for querying. To do this, first you must be logging to a CSV file, which here we will call pglog.csv.

First, install file_fdw as an extension:

```sql
CREATE EXTENSION file_fdw;
```

Then create a foreign server:

```sql
CREATE SERVER pglog FOREIGN DATA WRAPPER file_fdw;
```

Now you are ready to create the foreign data table. Using the CREATE FOREIGN TABLE command, you will need to define the columns for the table, the CSV file name, and its format:

```sql
CREATE FOREIGN TABLE pglog (  
    log_time timestamp(3) with time zone,  
    user_name text,  
```

3349
database_name text,
process_id integer,
connection_from text,
session_id text,
session_line_num bigint,
command_tag text,
session_start_time timestamp with time zone,
virtual_transaction_id text,
transaction_id bigint,
error_severity text,
sql_state_code text,
message text,
detail text,
hint text,
internal_query text,
internal_query_pos integer,
context text,
query text,
query_pos integer,
location text,
application_name text
) SERVER pglog
OPTIONS ( filename '/home/josh/data/log/pglog.csv', format 'csv' );

That's it — now you can query your log directly. In production, of course, you would need to define some way to deal with log rotation.

F.15. fuzzystrmatch

The fuzzystrmatch module provides several functions to determine similarities and distance between strings.

Caution

At present, the soundex, metaphone, dmetaphone, and dmetaphone_alt functions do not work well with multibyte encodings (such as UTF-8).

F.15.1. Soundex

The Soundex system is a method of matching similar-sounding names by converting them to the same code. It was initially used by the United States Census in 1880, 1900, and 1910. Note that Soundex is not very useful for non-English names.

The fuzzystrmatch module provides two functions for working with Soundex codes:

soundex(text) returns text
difference(text, text) returns int

The soundex function converts a string to its Soundex code. The difference function converts two strings to their Soundex codes and then reports the number of matching code positions. Since Soundex codes have four characters, the result ranges from zero to four, with zero being no match and four being an exact match. (Thus, the function is misnamed — similarity would have been a better name.)
Here are some usage examples:

```sql
SELECT soundex('hello world!');
SELECT soundex('Anne'), soundex('Ann'), difference('Anne', 'Ann');
SELECT soundex('Anne'), soundex('Andrew'), difference('Anne', 'Andrew');
SELECT soundex('Anne'), soundex('Margaret'), difference('Anne', 'Margaret');

CREATE TABLE s (nm text);
INSERT INTO s VALUES ('john');
INSERT INTO s VALUES ('joan');
INSERT INTO s VALUES ('wobbly');
INSERT INTO s VALUES ('jack');
SELECT * FROM s WHERE soundex(nm) = soundex('john');
SELECT * FROM s WHERE difference(s.nm, 'john') > 2;
```

### F.15.2. Levenshtein

This function calculates the Levenshtein distance between two strings:

```sql
levenshtein(text source, text target, int ins_cost, int del_cost, int sub_cost) returns int
levenshtein(text source, text target) returns int
levenshtein_less_equal(text source, text target, int ins_cost, int del_cost, int sub_cost, int max_d) returns int
levenshtein_less_equal(text source, text target, int max_d) returns int
```

Both source and target can be any non-null string, with a maximum of 255 characters. The cost parameters specify how much to charge for a character insertion, deletion, or substitution, respectively. You can omit the cost parameters, as in the second version of the function; in that case they all default to 1.

`levenshtein_less_equal` is an accelerated version of the Levenshtein function for use when only small distances are of interest. If the actual distance is less than or equal to `max_d`, then `levenshtein_less_equal` returns the correct distance; otherwise it returns some value greater than `max_d`. If `max_d` is negative then the behavior is the same as `levenshtein`.

Examples:

```sql
test=# SELECT levenshtein('GUMBO', 'GAMBOL');
levenshtein
----------
  2
(1 row)

test=# SELECT levenshtein('GUMBO', 'GAMBOL', 2, 1, 1);
levenshtein
----------
3351
```
F.15.3. Metaphone

Metaphone, like Soundex, is based on the idea of constructing a representative code for an input string. Two strings are then deemed similar if they have the same codes.

This function calculates the metaphone code of an input string:

```
metaphone(text source, int max_output_length) returns text
```

`source` has to be a non-null string with a maximum of 255 characters. `max_output_length` sets the maximum length of the output metaphone code; if longer, the output is truncated to this length.

Example:

```
test=# SELECT metaphone('GUMBO', 4);
metaphone
-----------
KM
   (1 row)
```

F.15.4. Double Metaphone

The Double Metaphone system computes two “sounds like” strings for a given input string — a “primary” and an “alternate”. In most cases they are the same, but for non-English names especially they can be a bit different, depending on pronunciation. These functions compute the primary and alternate codes:

```
dmetaphone(text source) returns text
dmetaphone_alt(text source) returns text
```

There is no length limit on the input strings.

Example:

```
test=# SELECT dmetaphone('gumbo');
dmetaphone
------------
3352
   (1 row)
```
F.16. hstore

This module implements the hstore data type for storing sets of key/value pairs within a single PostgreSQL value. This can be useful in various scenarios, such as rows with many attributes that are rarely examined, or semi-structured data. Keys and values are simply text strings.

F.16.1. hstore External Representation

The text representation of an hstore, used for input and output, includes zero or more key => value pairs separated by commas. Some examples:

k => v
foo => bar, baz => whatever
"1-a" => "anything at all"

The order of the pairs is not significant (and may not be reproduced on output). Whitespace between pairs or around the => sign is ignored. Double-quote keys and values that include whitespace, commas, =s or >s. To include a double quote or a backslash in a key or value, escape it with a backslash.

Each key in an hstore is unique. If you declare an hstore with duplicate keys, only one will be stored in the hstore and there is no guarantee as to which will be kept:

```sql
SELECT 'a=>1,a=>2'::hstore;

+-----------------+
| hstore          |
| "a"=>"1"        |
+-----------------+
```

A value (but not a key) can be an SQL NULL. For example:

```sql
key => NULL
```

The NULL keyword is case-insensitive. Double-quote the NULL to treat it as the ordinary string "NULL".

### Note

Keep in mind that the hstore text format, when used for input, applies before any required quoting or escaping. If you are passing an hstore literal via a parameter, then no additional processing is needed. But if you're passing it as a quoted literal constant, then any single-quote characters and (depending on the setting of the standard_conforming_strings configuration parameter) backslash characters need to be escaped correctly. See Section 4.1.2.1 for more on the handling of string constants.

On output, double quotes always surround keys and values, even when it's not strictly necessary.

F.16.2. hstore Operators and Functions

The operators provided by the hstore module are shown in Table F.7, the functions in Table F.8.
### Table F.7. hstore Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore -&gt; text</td>
<td>get value for key (NULL if not present)</td>
<td>'a'=&gt;x, b=&gt;y::hstore -&gt; 'a'</td>
<td></td>
</tr>
<tr>
<td>hstore -&gt; text[]</td>
<td>get values for keys (NULL if not present)</td>
<td>'a'=&gt;x, b=&gt;y, {&quot;z&quot;,&quot;x&quot;} hstore -&gt; c=x::hstore -&gt; ARRAY:['c','a']</td>
<td></td>
</tr>
<tr>
<td>hstore</td>
<td></td>
<td>hstore</td>
<td>concatenate hstores</td>
</tr>
<tr>
<td>hstore ? text</td>
<td>does hstore contain key?</td>
<td>'a'=&gt;1::hstore ? t 'a'</td>
<td></td>
</tr>
<tr>
<td>hstore ?&amp; text[]</td>
<td>does hstore contain all specified keys?</td>
<td>'a'=&gt;1,b=&gt;2::hstore ?&amp; ARRAY['a','b']</td>
<td></td>
</tr>
<tr>
<td>hstore ? text[]</td>
<td>does hstore contain any of the specified keys?</td>
<td>'a'=&gt;1,b=&gt;2::hstore ?</td>
<td>ARRAY['b','c']</td>
</tr>
<tr>
<td>hstore @&gt; hstore</td>
<td>does left operand contain right?</td>
<td>'a'=&gt;b, b=&gt;1, t c=&gt;NULL::hstore @&gt; 'b'=&gt;1</td>
<td></td>
</tr>
<tr>
<td>hstore @&gt; hstore</td>
<td>is left operand contained in right?</td>
<td>'a'=&gt;c::hstore f &lt;@ 'a'=&gt;b, b=&gt;1, c=&gt;NULL::hstore @&gt; 'b'=&gt;1</td>
<td></td>
</tr>
<tr>
<td>hstore - text</td>
<td>delete key from left operand</td>
<td>'a'=&gt;1, b=&gt;2, &quot;a&quot;=&gt;'1&quot;, c=&gt;3::hstore - &quot;c&quot;=&gt;'3&quot; 'b'::text</td>
<td></td>
</tr>
<tr>
<td>hstore - text[]</td>
<td>delete keys from left operand</td>
<td>'a'=&gt;1, b=&gt;2, &quot;c&quot;=&gt;'3&quot; c=&gt;3::hstore - ARRAY['a','b']</td>
<td></td>
</tr>
<tr>
<td>hstore - hstore</td>
<td>delete matching pairs from left operand</td>
<td>'a'=&gt;1, b=&gt;2, &quot;a&quot;=&gt;'1&quot;, c=&gt;3::hstore - ARRAY['a','b']</td>
<td></td>
</tr>
<tr>
<td>record #= hstore</td>
<td>replace fields in record with matching values from hstore</td>
<td>see Examples section</td>
<td></td>
</tr>
<tr>
<td>%% hstore</td>
<td>convert hstore to array of alternating keys and values</td>
<td>'a'=&gt;foo, {a,foo,b,bar} b=&gt;bar::hstore</td>
<td></td>
</tr>
<tr>
<td>%# hstore</td>
<td>convert hstore to two-dimensional key/value array</td>
<td>'a'=&gt;foo, {(a,foo), b=&gt;bar::hstore (b,bar)}</td>
<td></td>
</tr>
</tbody>
</table>

**Note**

Prior to PostgreSQL 8.2, the containment operators @> and <@ were called @ and ~, respectively. These names are still available, but are deprecated and will eventually be removed. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!
### Table F.8. hstore Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore(record)</td>
<td>hstore</td>
<td>construct an hstore from a record or row</td>
<td>hstore(ROW(1, 2))</td>
<td>f1=&gt;1, f2=&gt;2</td>
</tr>
<tr>
<td>hstore(tex-</td>
<td></td>
<td></td>
<td>hstore(ARRAY['a', 'b'])</td>
<td>a=&gt;1, b=&gt;2,</td>
</tr>
<tr>
<td>t[])</td>
<td>hstore</td>
<td>construct an hstore from an array, which may be a key/value array, or a two-dimensional array</td>
<td></td>
<td>hstore(ARRAY['c', '3'], ['d', '4'])</td>
</tr>
<tr>
<td>hstore(tex-</td>
<td></td>
<td></td>
<td>hstore(ARRAY['a', 'b'])</td>
<td>a=&gt;1, b=&gt;2,</td>
</tr>
<tr>
<td>t[], text[])</td>
<td></td>
<td>construct an hstore from separate key and value arrays</td>
<td></td>
<td>a=&gt;1, b=&gt;2,</td>
</tr>
<tr>
<td></td>
<td>hstore</td>
<td>make single-item hstore</td>
<td>hstore('a', 'b')</td>
<td>&quot;a&quot;&gt;=&quot;b&quot;</td>
</tr>
<tr>
<td>akeys(h-</td>
<td>text[]</td>
<td>get hstore's keys as an array</td>
<td>akeys('a=&gt;1,b=&gt;2')</td>
<td>{a,b}</td>
</tr>
<tr>
<td>store)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>skeys(h-</td>
<td>setof text</td>
<td>get hstore's keys as a set</td>
<td>skeys('a=&gt;1,b=&gt;2')</td>
<td>a,b</td>
</tr>
<tr>
<td>store)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>avals(h-</td>
<td>text[]</td>
<td>get hstore's values as an array</td>
<td>avals('a=&gt;1,b=&gt;2')</td>
<td>{1,2}</td>
</tr>
<tr>
<td>store)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>svals(h-</td>
<td>setof text</td>
<td>get hstore's values as a set</td>
<td>svals('a=&gt;1,b=&gt;2')</td>
<td></td>
</tr>
<tr>
<td>store)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>hstore_to_ar-</td>
<td>text[]</td>
<td>get hstore's keys and values as an array of alternating keys and values</td>
<td>hstore_to_ar-ray('a=&gt;1,b=&gt;2')</td>
<td>{a,1,b,2}</td>
</tr>
<tr>
<td>ray(hstore)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>hstore_to_ma-</td>
<td>text[]</td>
<td>get hstore's keys and values as a two-dimensional array</td>
<td>hstore_to_matrix('a=&gt;1,b=&gt;2')</td>
<td>2</td>
</tr>
<tr>
<td>trix(hstore)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
| | json | get hstore as a JSON value, converting all non-null values to JSON strings | hstore_to_json('"a key">="1", "b">="2"
| json_value, | | | | |
| | jsonb | get hstore as a JSONB value, converting all non-null values to JSON strings | hstore_to_jsonb('"a key">="1", "b">="2"
<p>| json_value, | | | | |
| | | | | |</p>
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore_to_json</td>
<td>json</td>
<td>get hstore as a json value, but attempt to distinguish numerical and Boolean values so they are unquoted in the JSON</td>
<td><code>&quot;a key&quot;=&gt;1, &quot;b&quot;=&gt;true, &quot;c&quot;=&gt;null, &quot;d&quot;=&gt;12345, &quot;e&quot;=&gt;&quot;012345&quot;, &quot;f&quot;=&gt;1.234, &quot;g&quot;=&gt;2.345e+4</code></td>
<td><code>{&quot;a key&quot;:1, &quot;b&quot;:true, &quot;c&quot;:null, &quot;d&quot;:12345, &quot;e&quot;:&quot;012345&quot;, &quot;f&quot;:1.234, &quot;g&quot;:2.345e+4}</code></td>
</tr>
<tr>
<td>son_loose(hstore)</td>
<td>jsonb</td>
<td>get hstore as a jsonb value, but attempt to distinguish numerical and Boolean values so they are unquoted in the JSON</td>
<td><code>&quot;a key&quot;=&gt;1, &quot;b&quot;=&gt;true, &quot;c&quot;=&gt;null, &quot;d&quot;=&gt;12345, &quot;e&quot;=&gt;&quot;012345&quot;, &quot;f&quot;=&gt;1.234, &quot;g&quot;=&gt;2.345e+4</code></td>
<td><code>{&quot;a key&quot;:1, &quot;b&quot;:true, &quot;c&quot;:null, &quot;d&quot;:12345, &quot;e&quot;:&quot;012345&quot;, &quot;f&quot;:1.234, &quot;g&quot;:2.345e+4}</code></td>
</tr>
<tr>
<td>slice(hstore, text[])</td>
<td>hstore</td>
<td>extract a subset of an hstore</td>
<td><code>slice('a=&gt;1,b=&gt;2,c=&gt;3'::hstore, ARRAY['b','c','x'])</code></td>
<td><code>{&quot;b&quot;:&quot;2&quot;,&quot;c&quot;:&quot;3&quot;}</code></td>
</tr>
<tr>
<td>each(hstore)</td>
<td>setof(key text, value text)</td>
<td>get hstore's keys and values as a set</td>
<td>select * from each('a=&gt;1,b=&gt;2')</td>
<td>`</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>------------</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>b</td>
</tr>
<tr>
<td>exist(hstore, text)</td>
<td>boolean</td>
<td>does hstore contain key?</td>
<td><code>exist('a=&gt;1','a')</code></td>
<td><code>true</code></td>
</tr>
<tr>
<td>defined(hstore, text)</td>
<td>boolean</td>
<td>does hstore contain non-NULL value for key?</td>
<td><code>defined('a=&gt;NUL-L','a')</code></td>
<td><code>false</code></td>
</tr>
<tr>
<td>delete(hstore, text)</td>
<td>hstore</td>
<td>delete pair with matching key</td>
<td><code>delete('a=&gt;1,b=&gt;2','b')</code></td>
<td><code>null</code></td>
</tr>
<tr>
<td>delete(hstore, text[])</td>
<td>hstore</td>
<td>delete pairs with matching keys</td>
<td><code>delete('a=&gt;1,b=&gt;2','c','x')</code></td>
<td><code>{&quot;c&quot;:&quot;3&quot;}</code></td>
</tr>
<tr>
<td>delete(hstore, hstore)</td>
<td>hstore</td>
<td>delete pairs matching those in the store</td>
<td><code>delete('a=&gt;1,b=&gt;2','a=&gt;4,b=&gt;2':hstore)</code></td>
<td><code>{&quot;a&quot;:&quot;1&quot;}</code></td>
</tr>
<tr>
<td>populate_record(record, hstore)</td>
<td></td>
<td>replace fields in see Examples section with values from hstore</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

3356
### F.16.3. Indexes

*hstore* has GiST and GIN index support for the `@>`, `?`, `?&` and `?|` operators. For example:

```sql
CREATE INDEX hidx ON testhstore USING GIST (h);
CREATE INDEX hidx ON testhstore USING GIN (h);
```

*hstore* also supports btree or hash indexes for the `=` operator. This allows *hstore* columns to be declared UNIQUE, or to be used in GROUP BY, ORDER BY or DISTINCT expressions. The sort ordering for *hstore* values is not particularly useful, but these indexes may be useful for equivalence lookups. Create indexes for `=` comparisons as follows:

```sql
CREATE INDEX hidx ON testhstore USING BTREE (h);
CREATE INDEX hidx ON testhstore USING HASH (h);
```

### F.16.4. Examples

Add a key, or update an existing key with a new value:

```sql
UPDATE tab SET h = h || hstore('c', '3');
```

Delete a key:

```sql
UPDATE tab SET h = delete(h, 'k1');
```

Convert a record to an *hstore*:

```sql
CREATE TABLE test (col1 integer, col2 text, col3 text);
INSERT INTO test VALUES (123, 'foo', 'bar');
SELECT hstore(t) FROM test AS t;
```

```
---------------------------------------------
"col1"=>"123", "col2"=>"foo", "col3"=>"bar"
(1 row)
```
Additional Supplied Modules

Convert an hstore to a predefined record type:

```
CREATE TABLE test (col1 integer, col2 text, col3 text);

SELECT * FROM populate_record(NULL::test,
    "col1"=>"456", "col2"=>"zzz");
```

<table>
<thead>
<tr>
<th>col1</th>
<th>col2</th>
<th>col3</th>
</tr>
</thead>
<tbody>
<tr>
<td>456</td>
<td>zzz</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>(1 row)</td>
</tr>
</tbody>
</table>

Modify an existing record using the values from an hstore:

```
CREATE TABLE test (col1 integer, col2 text, col3 text);
INSERT INTO test VALUES (123, 'foo', 'bar');

SELECT (r).* FROM (SELECT t #= "col3"=>'baz' AS r FROM test t) s;
```

<table>
<thead>
<tr>
<th>col1</th>
<th>col2</th>
<th>col3</th>
</tr>
</thead>
<tbody>
<tr>
<td>123</td>
<td>foo</td>
<td>baz</td>
</tr>
<tr>
<td></td>
<td></td>
<td>(1 row)</td>
</tr>
</tbody>
</table>

F.16.5. Statistics

The hstore type, because of its intrinsic liberality, could contain a lot of different keys. Checking for valid keys is the task of the application. The following examples demonstrate several techniques for checking keys and obtaining statistics.

Simple example:

```
SELECT * FROM each('aaa=>bq, b=>NULL, ""=>1');
```

Using a table:

```
SELECT (each(h)).key, (each(h)).value INTO stat FROM testhstore;
```

Online statistics:

```
SELECT key, count(*) FROM
(SELECT (each(h)).key FROM testhstore) AS stat
GROUP BY key
ORDER BY count DESC, key;
```

<table>
<thead>
<tr>
<th>key</th>
<th>count</th>
</tr>
</thead>
<tbody>
<tr>
<td>line</td>
<td>883</td>
</tr>
<tr>
<td>query</td>
<td>207</td>
</tr>
<tr>
<td>pos</td>
<td>203</td>
</tr>
<tr>
<td>node</td>
<td>202</td>
</tr>
<tr>
<td>space</td>
<td>197</td>
</tr>
<tr>
<td>status</td>
<td>195</td>
</tr>
<tr>
<td>public</td>
<td>194</td>
</tr>
<tr>
<td>title</td>
<td>190</td>
</tr>
<tr>
<td>org</td>
<td>189</td>
</tr>
<tr>
<td>.................</td>
<td>3358</td>
</tr>
</tbody>
</table>
F.16.6. Compatibility

As of PostgreSQL 9.0, hstore uses a different internal representation than previous versions. This presents no obstacle for dump/restore upgrades since the text representation (used in the dump) is unchanged.

In the event of a binary upgrade, upward compatibility is maintained by having the new code recognize old-format data. This will entail a slight performance penalty when processing data that has not yet been modified by the new code. It is possible to force an upgrade of all values in a table column by doing an UPDATE statement as follows:

```
UPDATE tablename SET hstorecol = hstorecol || '';
```

Another way to do it is:

```
ALTER TABLE tablename ALTER hstorecol TYPE hstore USING hstorecol || '';
```

The ALTER TABLE method requires an exclusive lock on the table, but does not result in bloating the table with old row versions.

F.16.7. Transforms

Additional extensions are available that implement transforms for the hstore type for the languages PL/Perl and PL/Python. The extensions for PL/Perl are called hstore_plperl and hstore_plperlu, for trusted and untrusted PL/Perl. If you install these transforms and specify them when creating a function, hstore values are mapped to Perl hashes. The extensions for PL/Python are called hstore_plpythonu, hstore_plpython2u, and hstore_plpython3u (see Section 46.1 for the PL/Python naming convention). If you use them, hstore values are mapped to Python dictionaries.

F.16.8. Authors

Oleg Bartunov <oleg@sai.msu.su>, Moscow, Moscow University, Russia

Teodor Sigaev <teodor@sigaev.ru>, Moscow, Delta-Soft Ltd., Russia

Additional enhancements by Andrew Gierth <andrew@taol1.riddles.org.uk>, United Kingdom

F.17. intagg

The intagg module provides an integer aggregator and an enumerator. intagg is now obsolete, because there are built-in functions that provide a superset of its capabilities. However, the module is still provided as a compatibility wrapper around the built-in functions.

F.17.1. Functions

The aggregator is an aggregate function `int_array_aggregate(integer)` that produces an integer array containing exactly the integers it is fed. This is a wrapper around `array_agg`, which does the same thing for any array type.

The enumerator is a function `int_array_enum(integer[])` that returns `setof integer`. It is essentially the reverse operation of the aggregator: given an array of integers, expand it into a set of rows. This is a wrapper around `unnest`, which does the same thing for any array type.
F.17.2. Sample Uses

Many database systems have the notion of a one to many table. Such a table usually sits between two indexed tables, for example:

```
CREATE TABLE left (id INT PRIMARY KEY, ...);
CREATE TABLE right (id INT PRIMARY KEY, ...);
CREATE TABLE one_to_many(left INT REFERENCES left, right INT REFERENCES right);
```

It is typically used like this:

```
SELECT right.* from right JOIN one_to_many ON (right.id = one_to_many.right)
  WHERE one_to_many.left = item;
```

This will return all the items in the right hand table for an entry in the left hand table. This is a very common construct in SQL.

Now, this methodology can be cumbersome with a very large number of entries in the `one_to_many` table. Often, a join like this would result in an index scan and a fetch for each right hand entry in the table for a particular left hand entry. If you have a very dynamic system, there is not much you can do. However, if you have some data which is fairly static, you can create a summary table with the aggregator.

```
CREATE TABLE summary AS
  SELECT left, int_array_aggregate(right) AS right
  FROM one_to_many
  GROUP BY left;
```

This will create a table with one row per left item, and an array of right items. Now this is pretty useless without some way of using the array; that's why there is an array enumerator. You can do

```
SELECT left, int_array_enum(right) FROM summary WHERE left = item;
```

The above query using `int_array_enum` produces the same results as

```
SELECT left, right FROM one_to_many WHERE left = item;
```

The difference is that the query against the summary table has to get only one row from the table, whereas the direct query against `one_to_many` must index scan and fetch a row for each entry.

On one system, an `EXPLAIN` showed a query with a cost of 8488 was reduced to a cost of 329. The original query was a join involving the `one_to_many` table, which was replaced by:

```
SELECT right, count(right) FROM
  ( SELECT left, int_array_enum(right) AS right
    FROM summary JOIN (SELECT left FROM left_table WHERE left = item) AS lefts
    ON (summary.left = lefts.left)
  ) AS list
GROUP BY right
ORDER BY count DESC;
```
F.18. intarray

The \texttt{intarray} module provides a number of useful functions and operators for manipulating null-free arrays of integers. There is also support for indexed searches using some of the operators.

All of these operations will throw an error if a supplied array contains any NULL elements.

Many of these operations are only sensible for one-dimensional arrays. Although they will accept input arrays of more dimensions, the data is treated as though it were a linear array in storage order.

F.18.1. \texttt{intarray} Functions and Operators

The functions provided by the \texttt{intarray} module are shown in Table F.9, the operators in Table F.10.

Table F.9. \texttt{intarray} Functions

\begin{tabular}{|l|l|l|l|}
\hline
Function & Return Type & Description & Example & Result \\
\hline
\texttt{icount(in-t[])} & \texttt{int} & number of elements in array & \texttt{icount('1,2,3':int[])} & 3 \\
\hline
\texttt{sort(int[], text dir)} & \texttt{int[]} & sort array — \texttt{dir} must be \texttt{asc} or \texttt{t[]}, \texttt{'desc'} & \texttt{sort('1,2,3':int[]), 'desc'} & desc \\
\hline
\texttt{sort(int[])} & \texttt{int[]} & sort in ascending order & \texttt{sort(array[11,44,77])} & (11,44,77) \\
\hline
\texttt{sort_asc(int[])} & \texttt{int[]} & sort in ascending order & \texttt{sort_asc(array[11,22,33,22,11])} & \\
\hline
\texttt{sort_desc(int[])} & \texttt{int[]} & sort in descending order & \texttt{sort_desc(array[11,22,33,22,11])} & \\
\hline
\texttt{uniq(int[])} & \texttt{int[]} & remove adjacent duplicates & \texttt{uniq(sort('1,2,3,2,3':int[]))} & \\
\hline
\texttt{idx(int[], int item)} & \texttt{int} & index of first element matching \texttt{item} & \texttt{idx(array[11,22,33,22,11], 22)} & 22 \\
\hline
\texttt{subarray(int[] t[], int start, int len)} & \texttt{int[]} & portion of array starting at position \texttt{start}, \texttt{len} elements & \texttt{subarray('1,2,3,2,1':int[], 2, 3)} & (2,3,2) \\
\hline
\texttt{intset(int)} & \texttt{int[]} & make single-element array & \texttt{intset(42)} & 42 \\
\hline
\end{tabular}

Table F.10. \texttt{intarray} Operators

\begin{tabular}{|l|l|l|}
\hline
Operator & Returns & Description \\
\hline
\texttt{int[] && int[]} & \texttt{boolean} & overlap — true if arrays have at least one common element \\
\hline
\texttt{int[] @> int[]} & \texttt{boolean} & contains — true if left array contains right array \\
\hline
\texttt{int[] <@ int[]} & \texttt{boolean} & contained — true if left array is contained in right array \\
\hline
\end{tabular}
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td># int[]</td>
<td>int</td>
<td>number of elements in array</td>
</tr>
<tr>
<td>int[] # int</td>
<td>int</td>
<td>index (same as idx function)</td>
</tr>
<tr>
<td>int[] + int</td>
<td>int[]</td>
<td>push element onto array (add it to end of array)</td>
</tr>
<tr>
<td>int[] + int[]</td>
<td>int[]</td>
<td>array concatenation (right array added to the end of left one)</td>
</tr>
<tr>
<td>int[] - int</td>
<td>int[]</td>
<td>remove entries matching right argument from array</td>
</tr>
<tr>
<td>int[] - int[]</td>
<td>int[]</td>
<td>remove elements of right array from left</td>
</tr>
<tr>
<td>int[]</td>
<td>int</td>
<td>union of arguments</td>
</tr>
<tr>
<td>int[]</td>
<td>int[]</td>
<td>union of arrays</td>
</tr>
<tr>
<td>int[] &amp; int[]</td>
<td>int[]</td>
<td>intersection of arrays</td>
</tr>
<tr>
<td>int[] @@ query_int</td>
<td>boolean</td>
<td>true if array satisfies query (see below)</td>
</tr>
<tr>
<td>query_int ~~ int[]</td>
<td>boolean</td>
<td>true if array satisfies query (comutator of @@)</td>
</tr>
</tbody>
</table>

(Before PostgreSQL 8.2, the containment operators @> and <@ were respectively called @ and ~. These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The operators &@, @>, and <@ are equivalent to PostgreSQL’s built-in operators of the same names, except that they work only on integer arrays that do not contain nulls, while the built-in operators work for any array type. This restriction makes them faster than the built-in operators in many cases.

The @@ and ~~ operators test whether an array satisfies a query, which is expressed as a value of a specialized data type query_int. A query consists of integer values that are checked against the elements of the array, possibly combined using the operators & (AND), | (OR), and ! (NOT). Parentheses can be used as needed. For example, the query 1& (2 | 3) matches arrays that contain 1 and also contain either 2 or 3.

#### F.18.2. Index Support

intarray provides index support for the &@, @>, <@, and @@ operators, as well as regular array equality.

Two GiST index operator classes are provided: gist__int_ops (used by default) is suitable for small- to medium-size data sets, while gist__intbig_ops uses a larger signature and is more suitable for indexing large data sets (i.e., columns containing a large number of distinct array values). The implementation uses an RD-tree data structure with built-in lossy compression.

There is also a non-default GIN operator class gin__int_ops supporting the same operators.

The choice between GiST and GIN indexing depends on the relative performance characteristics of GiST and GIN, which are discussed elsewhere.

#### F.18.3. Example

```sql
-- a message can be in one or more "sections"
CREATE TABLE message (mid INT PRIMARY KEY, sections INT[], ...);
```
Additional Supplied Modules

```
-- create specialized index
CREATE INDEX message_rdtree_idx ON message USING GIST (sections gist__int_ops);

-- select messages in section 1 OR 2 - OVERLAP operator
SELECT message.mid FROM message WHERE message.sections && '{1,2}';

-- select messages in sections 1 AND 2 - CONTAINS operator
SELECT message.mid FROM message WHERE message.sections @> '{1,2}';

-- the same, using QUERY operator
SELECT message.mid FROM message WHERE message.sections @@ '1&2'::query_int;
```

F.18.4. Benchmark

The source directory contrib/intarray/bench contains a benchmark test suite, which can be run against an installed PostgreSQL server. (It also requires DBD::Pg to be installed.) To run:

```bash
cd .../contrib/intarray/bench
createdb TEST
psql -c "CREATE EXTENSION intarray" TEST
./create_test.pl | psql TEST
./bench.pl
```

The bench.pl script has numerous options, which are displayed when it is run without any arguments.

F.18.5. Authors

All work was done by Teodor Sigaev (<teodor@sigaev.ru>) and Oleg Bartunov (<oleg@sai.msu.su>). See http://www.sai.msu.su/~megera/postgres/gist/ for additional information. Andrey Oktyabrski did a great work on adding new functions and operations.

F.19. isn

The isn module provides data types for the following international product numbering standards: EAN13, UPC, ISBN (books), ISMN (music), and ISSN (serials). Numbers are validated on input according to a hard-coded list of prefixes; this list of prefixes is also used to hyphenate numbers on output. Since new prefixes are assigned from time to time, the list of prefixes may be out of date. It is hoped that a future version of this module will obtained the prefix list from one or more tables that can be easily updated by users as needed; however, at present, the list can only be updated by modifying the source code and recompiling. Alternatively, prefix validation and hyphenation support may be dropped from a future version of this module.

F.19.1. Data Types

Table F.11 shows the data types provided by the isn module.

<table>
<thead>
<tr>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EAN13</td>
<td>European Article Numbers, always displayed in the EAN13 display format</td>
</tr>
</tbody>
</table>

3363
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ISBN13</td>
<td>International Standard Book Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISMN13</td>
<td>International Standard Music Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISSN13</td>
<td>International Standard Serial Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISBN</td>
<td>International Standard Book Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>ISMN</td>
<td>International Standard Music Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>ISSN</td>
<td>International Standard Serial Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>UPC</td>
<td>Universal Product Codes</td>
</tr>
</tbody>
</table>

Some notes:

1. ISBN13, ISMN13, ISSN13 numbers are all EAN13 numbers.
2. EAN13 numbers aren't always ISBN13, ISMN13 or ISSN13 (some are).
3. Some ISBN13 numbers can be displayed as ISBN.
4. Some ISMN13 numbers can be displayed as ISMN.
5. Some ISSN13 numbers can be displayed as ISSN.
6. UPC numbers are a subset of the EAN13 numbers (they are basically EAN13 without the first 0 digit).
7. All UPC, ISBN, ISMN and ISSN numbers can be represented as EAN13 numbers.

Internally, all these types use the same representation (a 64-bit integer), and all are interchangeable. Multiple types are provided to control display formatting and to permit tighter validity checking of input that is supposed to denote one particular type of number.

The ISBN, ISMN, and ISSN types will display the short version of the number (ISxN 10) whenever it's possible, and will show ISxN 13 format for numbers that do not fit in the short version. The EAN13, ISBN13, ISMN13 and ISSN13 types will always display the long version of the ISxN (EAN13).

### F.19.2. Casts

The isn module provides the following pairs of type casts:

- ISBN13 <> EAN13
- ISMN13 <> EAN13
- ISSN13 <> EAN13
- ISBN <> EAN13
- ISMN <> EAN13
- ISSN <> EAN13
- UPC <> EAN13

3364
Additional Supplied Modules

- ISMN <=> ISMN13
- ISSN <=> ISSN13

When casting from EAN13 to another type, there is a run-time check that the value is within the domain of the other type, and an error is thrown if not. The other casts are simply relabelings that will always succeed.

F.19.3. Functions and Operators

The isn module provides the standard comparison operators, plus B-tree and hash indexing support for all these data types. In addition there are several specialized functions; shown in Table F.12. In this table, isn means any one of the module’s data types.

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>isn_weak(boolean)</td>
<td>boolean</td>
<td>Sets the weak input mode (returns new setting)</td>
</tr>
<tr>
<td>isn_weak()</td>
<td>boolean</td>
<td>Gets the current status of the weak mode</td>
</tr>
<tr>
<td>make_valid(isn)</td>
<td>isn</td>
<td>Validates an invalid number (clears the invalid flag)</td>
</tr>
<tr>
<td>is_valid(isn)</td>
<td>boolean</td>
<td>Checks for the presence of the invalid flag</td>
</tr>
</tbody>
</table>

Weak mode is used to be able to insert invalid data into a table. Invalid means the check digit is wrong, not that there are missing numbers.

Why would you want to use the weak mode? Well, it could be that you have a huge collection of ISBN numbers, and that there are so many of them that for weird reasons some have the wrong check digit (perhaps the numbers were scanned from a printed list and the OCR got the numbers wrong, perhaps the numbers were manually captured... who knows). Anyway, the point is you might want to clean the mess up, but you still want to be able to have all the numbers in your database and maybe use an external tool to locate the invalid numbers in the database so you can verify the information and validate it more easily; so for example you’d want to select all the invalid numbers in the table.

When you insert invalid numbers in a table using the weak mode, the number will be inserted with the corrected check digit, but it will be displayed with an exclamation mark (!) at the end, for example 0-11-000322-5!. This invalid marker can be checked with the is_valid function and cleared with the make_valid function.

You can also force the insertion of invalid numbers even when not in the weak mode, by appending the ! character at the end of the number.

Another special feature is that during input, you can write ? in place of the check digit, and the correct check digit will be inserted automatically.

F.19.4. Examples

--Using the types directly:
SELECT isbn('978-0-393-04002-9');
SELECT isbn13('0901690546');
SELECT issn('1436-4522');

--Casting types:
Additional Supplied Modules

-- note that you can only cast from ean13 to another type when the
-- number would be valid in the realm of the target type;
-- thus, the following will NOT work: select
isbn(ean13('0220356483481'));
-- but these will:
SELECT upc(ean13('0220356483481'));
SELECT ean13(upc('220356483481'));

--Create a table with a single column to hold ISBN numbers:
CREATE TABLE test (id isbn);
INSERT INTO test VALUES('9780393040029');

--Automatically calculate check digits (observe the '?'):
INSERT INTO test VALUES('220500896?');
INSERT INTO test VALUES('978055215372?');
SELECT issn('3251231?');
SELECT ismn('979047213542?');

--Using the weak mode:
SELECT isn_weak(true);
INSERT INTO test VALUES('978-0-11-000533-4');
INSERT INTO test VALUES('9780141219307');
INSERT INTO test VALUES('2-205-00876-X');
SELECT isn_weak(false);

SELECT id FROM test WHERE NOT is_valid(id);
UPDATE test SET id = make_valid(id) WHERE id = '2-205-00876-X!';
SELECT * FROM test;
SELECT isbn13(id) FROM test;

F.19.5. Bibliography

The information to implement this module was collected from several sites, including:

- https://www.isbn-international.org/
- http://www.issn.org/
- https://www.ismn-international.org/
- https://www.wikipedia.org/

The prefixes used for hyphenation were also compiled from:

- https://www.gs1.org/standards/id-keys
- https://www.isbn-international.org/content/isbn-users-manual
- https://www.ismn-international.org/ranges.html

Care was taken during the creation of the algorithms and they were meticulously verified against the suggested algorithms in the official ISBN, ISMN, ISSN User Manuals.
F.19.6. Author

Germán Méndez Bravo (Kronuz), 2004 - 2006

This module was inspired by Garrett A. Wollman's isbn_issn code.

F.20. lo

The lo module provides support for managing Large Objects (also called LOs or BLOBs). This includes a data type lo and a trigger lo_manage.

F.20.1. Rationale

One of the problems with the JDBC driver (and this affects the ODBC driver also), is that the specification assumes that references to BLOBs (Binary Large OBjects) are stored within a table, and if that entry is changed, the associated BLOB is deleted from the database.

As PostgreSQL stands, this doesn't occur. Large objects are treated as objects in their own right; a table entry can reference a large object by OID, but there can be multiple table entries referencing the same large object OID, so the system doesn't delete the large object just because you change or remove one such entry.

Now this is fine for PostgreSQL-specific applications, but standard code using JDBC or ODBC won't delete the objects, resulting in orphan objects — objects that are not referenced by anything, and simply occupy disk space.

The lo module allows fixing this by attaching a trigger to tables that contain LO reference columns. The trigger essentially just does a lo_unlink whenever you delete or modify a value referencing a large object. When you use this trigger, you are assuming that there is only one database reference to any large object that is referenced in a trigger-controlled column!

The module also provides a data type lo, which is really just a domain of the oid type. This is useful for differentiating database columns that hold large object references from those that are OIDs of other things. You don't have to use the lo type to use the trigger, but it may be convenient to use it to keep track of which columns in your database represent large objects that you are managing with the trigger. It is also rumored that the ODBC driver gets confused if you don't use lo for BLOB columns.

F.20.2. How to Use It

Here's a simple example of usage:

CREATE TABLE image (title text, raster lo);

CREATE TRIGGER t_raster BEFORE UPDATE OR DELETE ON image
FOR EACH ROW EXECUTE FUNCTION lo_manage(raster);

For each column that will contain unique references to large objects, create a BEFORE UPDATE OR DELETE trigger, and give the column name as the sole trigger argument. You can also restrict the trigger to only execute on updates to the column by using BEFORE UPDATE OF column_name. If you need multiple lo columns in the same table, create a separate trigger for each one, remembering to give a different name to each trigger on the same table.

F.20.3. Limitations

- Dropping a table will still orphan any objects it contains, as the trigger is not executed. You can avoid this by preceding the DROP TABLE with DELETE FROM table.
TRUNCATE has the same hazard.

If you already have, or suspect you have, orphaned large objects, see the vacuumlo module to help you clean them up. It's a good idea to run vacuumlo occasionally as a back-stop to the lo_manage trigger.

• Some frontends may create their own tables, and will not create the associated trigger(s). Also, users may not remember (or know) to create the triggers.

F.20.4. Author

Peter Mount <peter@retep.org.uk>

F.21. ltree

This module implements a data type ltree for representing labels of data stored in a hierarchical tree-like structure. Extensive facilities for searching through label trees are provided.

F.21.1. Definitions

A label is a sequence of alphanumeric characters and underscores (for example, in C locale the characters A-Za-z0-9_ are allowed). Labels must be less than 256 bytes long.

Examples: 42, Personal_Services

A label path is a sequence of zero or more labels separated by dots, for example L1.L2.L3, representing a path from the root of a hierarchical tree to a particular node. The length of a label path must be less than 65kB, but keeping it under 2kB is preferable.

Example: Top.Countries.Europe.Russia

The ltree module provides several data types:

• ltree stores a label path.

• lquery represents a regular-expression-like pattern for matching ltree values. A simple word matches that label within a path. A star symbol (*) matches zero or more labels. For example:

  foo Match the exact label path foo
  *.foo.* Match any label path containing the label foo
  *.foo Match any label path whose last label is foo

Star symbols can also be quantified to restrict how many labels they can match:

  *{n} Match exactly n labels
  *(n,) Match at least n labels
  *(n,m) Match at least n but not more than m labels
  *{,m} Match at most m labels — same as *{0,m}

There are several modifiers that can be put at the end of a non-star label in lquery to make it match more than just the exact match:

  @ Match case-insensitively, for example a@ matches A
* Match any label with this prefix, for example foo*
matches foobar
%
Match initial underscore-separated words

The behavior of % is a bit complicated. It tries to match words rather than the entire label. For example foo_bar% matches foo_bar_baz but not foo_barbaz. If combined with *, prefix matching applies to each word separately, for example foo_bar%* matches fool_bar2_baz but not fool_br2_baz.

Also, you can write several possibly-modified labels separated with | (OR) to match any of those labels, and you can put ! (NOT) at the start to match any label that doesn't match any of the alternatives.

Here's an annotated example of lquery:

Top.*(0,2).sport@!.football|tennis.Russ*|Spain
a. b. c. d. e.

This query will match any label path that:

a. begins with the label Top
b. and next has zero to two labels before
c. a label beginning with the case-insensitive prefix sport
d. then a label not matching football nor tennis
e. and then ends with a label beginning with Russ or exactly matching Spain.

• ltxtquery represents a full-text-search-like pattern for matching ltree values. An ltxt-query value contains words, possibly with the modifiers @, *, % at the end; the modifiers have the same meanings as in lquery. Words can be combined with & (AND), | (OR), ! (NOT), and parentheses. The key difference from lquery is that ltxtquery matches words without regard to their position in the label path.

Here's an example ltxtquery:

Europe & Russia*@ & !Transportation

This will match paths that contain the label Europe and any label beginning with Russia (case-insensitive), but not paths containing the label Transportation. The location of these words within the path is not important. Also, when % is used, the word can be matched to any underscore-separated word within a label, regardless of position.

Note: ltxtquery allows whitespace between symbols, but ltree and lquery do not.

F.21.2. Operators and Functions

Type ltree has the usual comparison operators =, <>, <, <=, >=. Comparison sorts in the order of a tree traversal, with the children of a node sorted by label text. In addition, the specialized operators shown in Table F.13 are available.

Table F.13. ltree Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ltree @&gt; ltree</td>
<td>boolean</td>
<td>is left argument an ancestor of right (or equal)?</td>
</tr>
<tr>
<td>Operator</td>
<td>Returns</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>---------</td>
<td>--------------------------------------------------</td>
</tr>
<tr>
<td>ltree &lt;@ ltree</td>
<td>boolean</td>
<td>is left argument a descendant of right (or equal)?</td>
</tr>
<tr>
<td>lquery ~ ltree</td>
<td>boolean</td>
<td>does ltree match lquery?</td>
</tr>
<tr>
<td>ltree ? lquery[]</td>
<td>boolean</td>
<td>does ltree match any lquery in array?</td>
</tr>
<tr>
<td>lquery[] ? ltree</td>
<td>boolean</td>
<td>does ltree match any lquery in array?</td>
</tr>
<tr>
<td>ltree @ ltxtquery</td>
<td>boolean</td>
<td>does ltree match ltxt-query?</td>
</tr>
<tr>
<td>ltxtquery @ ltree</td>
<td>boolean</td>
<td>does ltree match ltxt-query?</td>
</tr>
<tr>
<td>ltree</td>
<td></td>
<td>ltree</td>
</tr>
<tr>
<td>ltree</td>
<td></td>
<td>text</td>
</tr>
<tr>
<td>text</td>
<td></td>
<td>ltree</td>
</tr>
<tr>
<td>ltree[] @&gt; ltree</td>
<td>boolean</td>
<td>does array contain an ancestor of ltree?</td>
</tr>
<tr>
<td>ltree &lt;@ ltree[]</td>
<td>boolean</td>
<td>does array contain an ancestor of ltree?</td>
</tr>
<tr>
<td>ltree[] &lt;@ ltree</td>
<td>boolean</td>
<td>does array contain a descendant of ltree?</td>
</tr>
<tr>
<td>ltree @&gt; ltree[]</td>
<td>boolean</td>
<td>does array contain a descendant of ltree?</td>
</tr>
<tr>
<td>ltree[] ~ lquery</td>
<td>boolean</td>
<td>does array contain any path matching lquery?</td>
</tr>
<tr>
<td>lquery ~ ltree[]</td>
<td>boolean</td>
<td>does array contain any path matching lquery?</td>
</tr>
<tr>
<td>ltree[] ? lquery[]</td>
<td>boolean</td>
<td>does ltree array contain any path matching any lquery?</td>
</tr>
<tr>
<td>lquery[] ? ltree[]</td>
<td>boolean</td>
<td>does ltree array contain any path matching any lquery?</td>
</tr>
<tr>
<td>ltree[] @ ltxtquery</td>
<td>boolean</td>
<td>does array contain any path matching ltxtquery?</td>
</tr>
<tr>
<td>ltxtquery @ ltree[]</td>
<td>boolean</td>
<td>does array contain any path matching ltxtquery?</td>
</tr>
<tr>
<td>ltree[] ?@&gt; ltree</td>
<td>ltree</td>
<td>first array entry that is an ancestor of ltree: NULL if none</td>
</tr>
<tr>
<td>ltree[] ?&lt;@ ltree</td>
<td>ltree</td>
<td>first array entry that is a descendant of ltree: NULL if none</td>
</tr>
<tr>
<td>ltree[] ?~ lquery</td>
<td>ltree</td>
<td>first array entry that matches lquery: NULL if none</td>
</tr>
<tr>
<td>ltree[] ?@ ltxtquery</td>
<td>ltree</td>
<td>first array entry that matches ltxtquery: NULL if none</td>
</tr>
</tbody>
</table>

The operators <@, @>, @ and ~ have analogues ^<@, ^@, ^@, ^~, which are the same except they do not use indexes. These are useful only for testing purposes.
The available functions are shown in Table F.14.

### Table F.14. ltree Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>subltree(ltree, int start, int end)</code></td>
<td><code>ltree</code></td>
<td>subpath of <code>ltree</code> starting from position <code>start</code> to position <code>end-1</code> (counting from 0)</td>
<td><code>subltree('Top.Child1.Child2',1,2)</code></td>
<td><code>Child1.Child2</code></td>
</tr>
<tr>
<td><code>subpath(ltree, int offset, int len)</code></td>
<td><code>ltree</code></td>
<td>subpath of <code>ltree</code> starting at position <code>offset</code>, length <code>len</code>. If <code>offset</code> is negative, subpath starts that far from the end of the path. If <code>len</code> is negative, leaves that many labels off the end of the path.</td>
<td><code>subpath('Top.Child1.Child2',0,2)</code></td>
<td><code>Top.Child1</code></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>nlevel(ltree)</code></td>
<td><code>integer</code></td>
<td>number of labels in path</td>
<td><code>nlevel('Top.Child1.Child2')</code></td>
<td><code>3</code></td>
</tr>
<tr>
<td><code>index(ltree a, ltree b)</code></td>
<td><code>integer</code></td>
<td>position of first occurrence of <code>b</code> in <code>a</code>; <code>dex('0.1.2.3.5.4.5.6.8.5.6.8','5.6')</code> if not found</td>
<td><code>index('0.1.2.3.5.4.5.6.8.5.6.8','5.6')</code></td>
<td><code>6</code></td>
</tr>
<tr>
<td><code>index(ltree a, ltree b, int offset)</code></td>
<td><code>integer</code></td>
<td>position of first occurrence of <code>b</code> in <code>a</code>, searching starting at <code>offset</code>; negative <code>offset</code> means start <code>offset</code> labels from the end of the path</td>
<td><code>index('0.1.2.3.5.4.5.6.8.5.6.8','5.6',-4)</code></td>
<td><code>9</code></td>
</tr>
<tr>
<td><code>text2ltree(text)</code></td>
<td><code>ltree</code></td>
<td>cast text to <code>ltree</code></td>
<td><code>text2ltree('text')</code></td>
<td><code>ltree</code></td>
</tr>
<tr>
<td><code>ltree2txt(ltree)</code></td>
<td><code>text</code></td>
<td>cast <code>ltree</code> to <code>text</code></td>
<td><code>ltree2txt('ltree')</code></td>
<td><code>text</code></td>
</tr>
<tr>
<td><code>lca(ltree, ...)</code></td>
<td><code>ltree</code></td>
<td>longest common ancestor of paths (up to 8 arguments supported)</td>
<td><code>lca('1.2.3', '1.2.3.4.5.6')</code></td>
<td><code>1.2</code></td>
</tr>
<tr>
<td><code>lca(ltree[])</code></td>
<td><code>ltree</code></td>
<td>longest common ancestor of paths in <code>array[1.2.3::ltree,'1.2.3.4']</code></td>
<td><code>lca(ltree[])</code></td>
<td><code>1.2</code></td>
</tr>
</tbody>
</table>
F.21.3. Indexes

ltree supports several types of indexes that can speed up the indicated operators:

- B-tree index over ltree: <, <=, =, >=, >
- GiST index over ltree: <, <=, =, >=, @>, <@, @, ~, ?

Example of creating such an index:

```
CREATE INDEX path_gist_idx ON test USING GIST (path);
```

Example of creating such an index:

```
CREATE INDEX path_gist_idx ON test USING GIST (array_path);
```

Note: This index type is lossy.

F.21.4. Example

This example uses the following data (also available in file contrib/ltree/ltreetest.sql in the source distribution):

```
CREATE TABLE test (path ltree);
INSERT INTO test VALUES ('Top');
INSERT INTO test VALUES ('Top.Science');
INSERT INTO test VALUES ('Top.Science.Astronomy');
INSERT INTO test VALUES ('Top.Science.Astronomy.Astrophysics');
INSERT INTO test VALUES ('Top.Science.Astronomy.Cosmology');
INSERT INTO test VALUES ('Top.Hobbies');
INSERT INTO test VALUES ('Top.Hobbies.Amateurs_Astronomy');
INSERT INTO test VALUES ('Top.Collections');
INSERT INTO test VALUES ('Top.Collections.Pictures');
INSERT INTO test VALUES ('Top.Collections.Pictures.Astronomy');
```

```
CREATE INDEX path_gist_idx ON test USING GIST (path);
CREATE INDEX path_idx ON test USING BTREE (path);
```

Now, we have a table test populated with data describing the hierarchy shown below:

```
Top
 / | | |
Science Hobbies Collections
 / | | |
Astronomy Amateurs_Astronomy Pictures
 / | |
Astrophysics Cosmology Astronomy
```

3372
We can do inheritance:

```
\$treetest=>\ SELECT\ path\ FROM\ test\ WHERE\ path\ <@\ 'Top.Science';

path
--------------
Top.Science
Top.Science.Astronomy
Top.Science.Astronomy.Astrophysics
(4\ rows)
```

Here are some examples of path matching:

```
\$treetest=>\ SELECT\ path\ FROM\ test\ WHERE\ path\ ~\ '*.Astronomy.*';

path
-----------------------------------------------
Top.Science.Astronomy
Top.Science.Astronomy.Astrophysics
Top.Collections.Pictures.Astronomy
(7\ rows)

\$treetest=>\ SELECT\ path\ FROM\ test\ WHERE\ path\ ~\ '*.!\ pictures@.*.Astronomy.*';

path
--------------
Top.Science.Astronomy
Top.Science.Astronomy.Astrophysics
(3\ rows)
```

Here are some examples of full text search:

```
\$treetest=>\ SELECT\ path\ FROM\ test\ WHERE\ path\ @\ 'Astro*\&\ !\ pictures@';

path
--------------
Top.Science.Astronomy
Top.Science.Astronomy.Astrophysics
Top.Hobbies.Amateurs_Astronomy
(4\ rows)

\$treetest=>\ SELECT\ path\ FROM\ test\ WHERE\ path\ @\ 'Astro*\&\ !\ pictures@';

path
--------------
Top.Science.Astronomy
Top.Science.Astronomy.Astrophysics
(3\ rows)
```
Path construction using functions:

```sql
ltreetest=> SELECT subpath(path,0,2)||'Space'||subpath(path,2) FROM test WHERE path <@ 'Top.Science.Astronomy';
?column?
------------------------------------------------------------
Top.Science.Space.Astronomy
(3 rows)
```

We could simplify this by creating a SQL function that inserts a label at a specified position in a path:

```sql
CREATE FUNCTION ins_label(ltree, int, text) RETURNS ltree
AS 'select subpath($1,0,$2) || $3 || subpath($1,$2);'
LANGUAGE SQL IMMUTABLE;

ltreetest=> SELECT ins_label(path,2,'Space') FROM test WHERE path <@ 'Top.Science.Astronomy';
ins_label
------------------------------------------------------------
Top.Science.Space.Astronomy
(3 rows)
```

**F.21.5. Transforms**

Additional extensions are available that implement transforms for the `ltree` type for PL/Python. The extensions are called `ltree_plpythonu`, `ltree_plpython2u`, and `ltree_plpython3u` (see Section 46.1 for the PL/Python naming convention). If you install these transforms and specify them when creating a function, `ltree` values are mapped to Python lists. (The reverse is currently not supported, however.)

**F.21.6. Authors**

All work was done by Teodor Sigaev (<teodor@stack.net>) and Oleg Bartunov (<oleg@sai.msu.su>). See http://www.sai.msu.su/~megera/postgres/gist/ for additional information. Authors would like to thank Eugeny Rodichev for helpful discussions. Comments and bug reports are welcome.

**F.22. pageinspect**

The `pageinspect` module provides functions that allow you to inspect the contents of database pages at a low level, which is useful for debugging purposes. All of these functions may be used only by superusers.

**F.22.1. General Functions**

```sql
get_raw_page(relname text, fork text, blkno int) returns bytea

get_raw_page reads the specified block of the named relation and returns a copy as a bytea value. This allows a single time-consistent copy of the block to be obtained. fork should be 'main' for the main data fork, 'fsm' for the free space map, 'vm' for the visibility map, or 'init' for the initialization fork.
```
get_raw_page(relname text, blkno int) returns bytea

A shorthand version of get_raw_page, for reading from the main fork. Equivalent to
get_raw_page(relname, 'main', blkno)

page_header(page bytea) returns record

page_header shows fields that are common to all PostgreSQL heap and index pages.

A page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT * FROM page_header(get_raw_page('pg_class', 0));

<table>
<thead>
<tr>
<th>lsn</th>
<th>checksum</th>
<th>flags</th>
<th>lower</th>
<th>upper</th>
<th>special</th>
<th>pagesize</th>
<th>version</th>
<th>prune_xid</th>
</tr>
</thead>
<tbody>
<tr>
<td>0/24A1B50</td>
<td>0</td>
<td>1</td>
<td>232</td>
<td>368</td>
<td>8192</td>
<td>8192</td>
<td>4</td>
<td>0</td>
</tr>
</tbody>
</table>

The returned columns correspond to the fields in the PageHeaderData struct. See src/ind-
clude/storage/bufpage.h for details.

The checksum field is the checksum stored in the page, which might be incorrect if the page
is somehow corrupted. If data checksums are not enabled for this instance, then the value stored
is meaningless.

page_checksum(page bytea, blkno int4) returns smallint

page_checksum computes the checksum for the page, as if it was located at the given block.

A page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT page_checksum(get_raw_page('pg_class', 0), 0);

<table>
<thead>
<tr>
<th>page_checksum</th>
</tr>
</thead>
<tbody>
<tr>
<td>13443</td>
</tr>
</tbody>
</table>

Note that the checksum depends on the block number, so matching block numbers should be
passed (except when doing esoteric debugging).

The checksum computed with this function can be compared with the checksum result field of
the function page_header. If data checksums are enabled for this instance, then the two values
should be equal.

heap_page_items(page bytea) returns setof record

heap_page_items shows all line pointers on a heap page. For those line pointers that are in
use, tuple headers as well as tuple raw data are also shown. All tuples are shown, whether or not
the tuples were visible to an MVCC snapshot at the time the raw page was copied.

A heap page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT * FROM heap_page_items(get_raw_page('pg_class', 0));

See src/include/storage/itemid.h and src/include/access/htup_de-
tails.h for explanations of the fields returned.
tuple_data_split(rel_oid, t_data bytea, t_infomask integer, t_infomask2 integer, t_bits text [, do_detoast bool]) returns bytea[]

tuple_data_split splits tuple data into attributes in the same way as backend internals.

test=# SELECT tuple_data_split('pg_class'::regclass, 
    t_data, t_infomask, t_infomask2, t_bits) FROM 
    heap_page_items(get_raw_page('pg_class', 0));

This function should be called with the same arguments as the return attributes of 
heap_page_items.

If do_detoast is true, attribute that will be detoasted as needed. Default value is false.

heap_page_item_attrs(rel_oid, t_data bytea, [, do_detoast bool]) returns bytea[]

heap_page_item_attrs is equivalent to heap_page_items except that it returns tuple 
raw data as an array of attributes that can optionally be detoasted by do_detoast which is 
false by default.

A heap page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT * FROM 
    heap_page_item_attrs(get_raw_page('pg_class', 0), 
    'pg_class'::regclass);

fsm_page_contents(page bytea) returns text

fsm_page_contents shows the internal node structure of a FSM page. The output is a mul-
tiline string, with one line per node in the binary tree within the page. Only those nodes that are 
not zero are printed. The so-called "next" pointer, which points to the next slot to be returned 
from the page, is also printed.

See src/backend/storage/freespace/README for more information on the structure 
of an FSM page.

F.22.2. B-tree Functions

bt_metap(relname text) returns record

bt_metap returns information about a B-tree index's metapage. For example:

test=# SELECT * FROM bt_metap('pg_cast_oid_index');

-[- RECORD 1 ]-----------------------
magic                | 340322
version              | 3
root                 | 1
level                | 0
fastroot             | 1
fastlevel            | 0
oldest_xact          | 582
last_cleanup_num_tuples | 1000

bt_page_stats(relname text, blkno int) returns record

bt_page_stats returns summary information about single pages of B-tree indexes. For ex-
ample:
test=# SELECT * FROM bt_page_stats('pg_cast_oid_index', 1);
-[ RECORD 1 ]----------
| blkno  | 1 |
| type   | 1 |
| live_items | 256 |
| dead_items | 0 |
| avg_item_size | 12 |
| page_size | 8192 |
| free_size | 4056 |
| btpo_prev | 0 |
| btpo_next | 0 |
| btpo | 0 |
| btpo_flags | 3 |

bt_page_items(relname text, blkno int) returns setof record

bt_page_items returns detailed information about all of the items on a B-tree index page. For example:

```sql
test=# SELECT * FROM bt_page_items('pg_cast_oid_index', 1);
```
```
| itemoffset | ctid   | itemlen | nulls | vars | data     |
|------------+--------+---------+-------+------+----------|
| 1          | (0,1)  | 12      | f     | f    | 23 27 00 00 |
| 2          | (0,2)  | 12      | f     | f    | 24 27 00 00 |
| 3          | (0,3)  | 12      | f     | f    | 25 27 00 00 |
| 4          | (0,4)  | 12      | f     | f    | 26 27 00 00 |
| 5          | (0,5)  | 12      | f     | f    | 27 27 00 00 |
| 6          | (0,6)  | 12      | f     | f    | 28 27 00 00 |
| 7          | (0,7)  | 12      | f     | f    | 29 27 00 00 |
| 8          | (0,8)  | 12      | f     | f    | 2a 27 00 00 |
```

In a B-tree leaf page, ctid points to a heap tuple. In an internal page, the block number part of ctid points to another page in the index itself, while the offset part (the second number) is ignored and is usually 1.

Note that the first item on any non-rightmost page (any page with a non-zero value in the btpo_next field) is the page’s “high key”, meaning its data serves as an upper bound on all items appearing on the page, while its ctid field is meaningless. Also, on non-leaf pages, the first real data item (the first item that is not a high key) is a “minus infinity” item, with no actual value in its data field. Such an item does have a valid downlink in its ctid field, however.

bt_page_items(page bytea) returns setof record

It is also possible to pass a page to bt_page_items as a bytea value. A page image obtained with get_raw_page should be passed as argument. So the last example could also be rewritten like this:

```sql
test=# SELECT * FROM
    bt_page_items(get_raw_page('pg_cast_oid_index', 1));
```
```
| itemoffset | ctid   | itemlen | nulls | vars | data     |
|------------+--------+---------+-------+------+----------|
| 1          | (0,1)  | 12      | f     | f    | 23 27 00 00 |
| 2          | (0,2)  | 12      | f     | f    | 24 27 00 00 |
| 3          | (0,3)  | 12      | f     | f    | 25 27 00 00 |
| 4          | (0,4)  | 12      | f     | f    | 26 27 00 00 |
| 5          | (0,5)  | 12      | f     | f    | 27 27 00 00 |
| 6          | (0,6)  | 12      | f     | f    | 28 27 00 00 |
```
F.22.3. BRIN Functions

brin_page_type(page bytea) returns text

brin_page_type returns the page type of the given BRIN index page, or throws an error if the page is not a valid BRIN page. For example:

test=# SELECT brin_page_type(get_raw_page('brinidx', 0));

brin_metapage_info(page bytea) returns record

brin_metapage_info returns assorted information about a BRIN index metapage. For example:

test=# SELECT * FROM brin_metapage_info(get_raw_page('brinidx', 0));

brin_revmap_data(page bytea) returns setof tid

brin_revmap_data returns the list of tuple identifiers in a BRIN index range map page. For example:

test=# SELECT * FROM brin_revmap_data(get_raw_page('brinidx', 2)) LIMIT 5;

brin_page_items(page bytea, index oid) returns setof record

brin_page_items returns the data stored in the BRIN data page. For example:

test=# SELECT * FROM brin_page_items(get_raw_page('brinidx', 5),
bronixd'", 0));

---------------
| magic    | version | pagesperrange | lastrevmappage |
|-----------+---------+---------------+----------------|
| 0xA8109CFA | 1       | 4             | 2              |

---------------
<table>
<thead>
<tr>
<th>pages</th>
</tr>
</thead>
<tbody>
<tr>
<td>(6,137)</td>
</tr>
<tr>
<td>(6,138)</td>
</tr>
<tr>
<td>(6,139)</td>
</tr>
<tr>
<td>(6,140)</td>
</tr>
<tr>
<td>(6,141)</td>
</tr>
</tbody>
</table>

---------------
| itemoffset  | blknum | attnum | allnulls | hasnulls |
|-------------+--------+--------+----------+----------|
| +------------+--------+--------+----------+----------|
| 137 | 0 | 1 | t | f | f |
The returned columns correspond to the fields in the BrinMemTuple and BrinValues structs. See src/include/access/brin_tuple.h for details.

F.22.4. GIN Functions

gin_metapage_info(page bytea) returns record

gin_metapage_info returns information about a GIN index metapage. For example:

test=# SELECT * FROM gin_metapage_info(get_raw_page('gin_index', 0));
-| [ RECORD 1 ]--------
  | pending_head | 4294967295
  | pending_tail | 4294967295
  | tail_free_size | 0
  | n_pending_pages | 0
  | n_pending_tuples | 0
  | n_total_pages | 7
  | n_entry_pages | 6
  | n_data_pages | 0
  | n_entries | 693
  | version | 2


gin_page_opaque_info(page bytea) returns record

 gin_page_opaque_info returns information about a GIN index opaque area, like the page type. For example:

 test=# SELECT * FROM gin_page_opaque_info(get_raw_page('gin_index', 2));
 rightlink | maxoff | flags
 ----------+--------+------------------------
  5 | 0 | {data,leaf,compressed}
(1 row)


gin_leafpage_items(page bytea) returns setof record

 gin_leafpage_items returns information about the data stored in a GIN leaf page. For example:

 test=# SELECT first_tid, nbytes, tids[0:5] AS some_tids
 FROM gin_leafpage_items(get_raw_page('gin_test_idx', 2));
 first_tid | nbytes | some_tids
 ----------|--------+--------------------------
 +--------------------------
F.22.5. Hash Functions

hash_page_type(page bytea) returns text

hash_page_type returns page type of the given HASH index page. For example:

test=# SELECT hash_page_type(get_raw_page('con_hash_index', 0));
    hash_page_type
  ----------------
     metapage

hash_page_stats(page bytea) returns setof record

hash_page_stats returns information about a bucket or overflow page of a HASH index. For example:

test=# SELECT * FROM
    hash_page_stats(get_raw_page('con_hash_index', 1));
          -[ RECORD 1 ]--
    live_items | 407
    dead_items |    0
    page_size  | 8192
    free_size  |    8
    hasho_prevblkno | 4096
    hasho_nextblkno | 8474
    hasho_bucket |    0
    hasho_flag  |   66
    hasho_page_id | 65408

hash_page_items(page bytea) returns setof record

hash_page_items returns information about the data stored in a bucket or overflow page of a HASH index page. For example:

test=# SELECT * FROM
    hash_page_items(get_raw_page('con_hash_index', 1)) LIMIT 5;
          itemoffset |   ctid    |    data
         ------------+-----------+------------
             1 | (899,77)  | 1053474816
             2 | (897,29)  | 1053474816
             3 | (894,207) | 1053474816
hash_bitmap_info(index oid, blkno int) returns record

hash_bitmap_info shows the status of a bit in the bitmap page for a particular overflow page of HASH index. For example:

```
test=# SELECT * FROM hash_bitmap_info('con_hash_index', 2052);
bitmapblkno | bitmapbit | bitstatus
-------------+-----------+-----------
          65 |          3 |         t
```

hash_metapage_info(page bytea) returns record

hash_metapage_info returns information stored in meta page of a HASH index. For example:

```
test=# SELECT magic, version, ntuples, ffactor, bsize, bmsize, bmshift,
      maxbucket, highmask, lowmask, ovflpoint, firstfree, nmaps, procid,
      regexp_replace(spares::text, '(,0)*}', '}') as spares,
      regexp_replace(mapp::text, '(,0)*}', '}') as mapp
      FROM hash_metapage_info(get_raw_page('con_hash_index', 0));
```

F.23. passwordcheck

The passwordcheck module checks users' passwords whenever they are set with CREATE ROLE or ALTER ROLE. If a password is considered too weak, it will be rejected and the command will terminate with an error.

To enable this module, add '$libdir/passwordcheck' to shared_preload_libraries in postgresql.conf, then restart the server.
You can adapt this module to your needs by changing the source code. For example, you can use CrackLib\(^2\) to check passwords — this only requires uncommenting two lines in the `Makefile` and rebuilding the module. (We cannot include CrackLib by default for license reasons.) Without CrackLib, the module enforces a few simple rules for password strength, which you can modify or extend as you see fit.

**Caution**

To prevent unencrypted passwords from being sent across the network, written to the server log or otherwise stolen by a database administrator, PostgreSQL allows the user to supply pre-encrypted passwords. Many client programs make use of this functionality and encrypt the password before sending it to the server.

This limits the usefulness of the `passwordcheck` module, because in that case it can only try to guess the password. For this reason, `passwordcheck` is not recommended if your security requirements are high. It is more secure to use an external authentication method such as GSSAPI (see Chapter 20) than to rely on passwords within the database.

Alternatively, you could modify `passwordcheck` to reject pre-encrypted passwords, but forcing users to set their passwords in clear text carries its own security risks.

### F.24. pg_buffercache

The `pg_buffercache` module provides a means for examining what's happening in the shared buffer cache in real time.

The module provides a C function `pg_buffercache_pages` that returns a set of records, plus a view `pg_buffercache` that wraps the function for convenient use.

By default use is restricted to superusers and members of the `pg_read_all_stats` role. Access may be granted to others using `GRANT`.

#### F.24.1. The pg_buffercache View

The definitions of the columns exposed by the view are shown in Table F.15.

**Table F.15. pg_buffercache Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bufferid</td>
<td>integer</td>
<td></td>
<td>ID, in the range 1..shared_buffers</td>
</tr>
<tr>
<td>relfilenode</td>
<td>oid</td>
<td><code>pg_class.relfilenode</code></td>
<td>Filenode number of the relation</td>
</tr>
<tr>
<td>reltablespace</td>
<td>oid</td>
<td><code>pg_table-space.oid</code></td>
<td>Tablespace OID of the relation</td>
</tr>
<tr>
<td>reldatabase</td>
<td>oid</td>
<td><code>pg_database.oid</code></td>
<td>Database OID of the relation</td>
</tr>
<tr>
<td>relforknumber</td>
<td>smallint</td>
<td></td>
<td>Fork number within the relation</td>
</tr>
</tbody>
</table>

\(^2\) https://sourceforge.net/projects/cracklib/
There is one row for each buffer in the shared cache. Unused buffers are shown with all fields null except bufferid. Shared system catalogs are shown as belonging to database zero.

Because the cache is shared by all the databases, there will normally be pages from relations not belonging to the current database. This means that there may not be matching join rows in pg_class for some rows, or that there could even be incorrect joins. If you are trying to join against pg_class, it’s a good idea to restrict the join to rows having reldatabase equal to the current database’s OID or zero.

When the pg_buffercache view is accessed, internal buffer manager locks are taken for long enough to copy all the buffer state data that the view will display. This ensures that the view produces a consistent set of results, while not blocking normal buffer activity longer than necessary. Nonetheless there could be some impact on database performance if this view is read often.

### F.24.2. Sample Output

```sql
regression=# SELECT c.relname, count(*) AS buffers
    FROM pg_buffercache b INNER JOIN pg_class c
    ON b.relfilenode = pg_relation_filenode(c.oid) AND
    b.reldatabase IN (0, (SELECT oid FROM pg_database
    WHERE datname = current_database()))
    GROUP BY c.relname
    ORDER BY 2 DESC
    LIMIT 10;
```

<table>
<thead>
<tr>
<th>relname</th>
<th>buffers</th>
</tr>
</thead>
<tbody>
<tr>
<td>tenk2</td>
<td>345</td>
</tr>
<tr>
<td>tenk1</td>
<td>141</td>
</tr>
<tr>
<td>pg_proc</td>
<td>46</td>
</tr>
<tr>
<td>pg_class</td>
<td>45</td>
</tr>
<tr>
<td>pg_attribute</td>
<td>43</td>
</tr>
<tr>
<td>pg_class_relnam_nsp_index</td>
<td>30</td>
</tr>
<tr>
<td>pg_proc_proname_args_nsp_index</td>
<td>28</td>
</tr>
<tr>
<td>pg_attribute_reld_attnam_index</td>
<td>26</td>
</tr>
<tr>
<td>pg_depend</td>
<td>22</td>
</tr>
<tr>
<td>pg_depend_reference_index</td>
<td>20</td>
</tr>
</tbody>
</table>

(10 rows)

### F.24.3. Authors

Mark Kirkwood <markir@paradise.net.nz>
Design suggestions: Neil Conway <neilc@samurai.com>
Debugging advice: Tom Lane <tgl@sss.pgh.pa.us>

F.25. pgcrypto

The `pgcrypto` module provides cryptographic functions for PostgreSQL.

F.25.1. General Hashing Functions

F.25.1.1. `digest()`

digest(data text, type text) returns bytea  
digest(data bytea, type text) returns bytea

Computes a binary hash of the given `data`. `type` is the algorithm to use. Standard algorithms are `md5`, `sha1`, `sha224`, `sha256`, `sha384` and `sha512`. If `pgcrypto` was built with OpenSSL, more algorithms are available, as detailed in Table F.19.

If you want the digest as a hexadecimal string, use `encode()` on the result. For example:

```sql
CREATE OR REPLACE FUNCTION sha1(bytea) returns text AS $$
SELECT encode(digest($1, 'sha1'), 'hex')
$$ LANGUAGE SQL STRICT IMMUTABLE;
```

F.25.1.2. `hmac()`

hmac(data text, key text, type text) returns bytea  
hmac(data bytea, key bytea, type text) returns bytea

Calculates hashed MAC for `data` with key `key`. `type` is the same as in `digest()`.

This is similar to `digest()` but the hash can only be recalculated knowing the key. This prevents the scenario of someone altering data and also changing the hash to match.

If the key is larger than the hash block size it will first be hashed and the result will be used as key.

F.25.2. Password Hashing Functions

The functions `crypt()` and `gen_salt()` are specifically designed for hashing passwords. `crypt()` does the hashing and `gen_salt()` prepares algorithm parameters for it.

The algorithms in `crypt()` differ from the usual MD5 or SHA1 hashing algorithms in the following respects:

1. They are slow. As the amount of data is so small, this is the only way to make brute-forcing passwords hard.

2. They use a random value, called the `salt`, so that users having the same password will have different encrypted passwords. This is also an additional defense against reversing the algorithm.

3. They include the algorithm type in the result, so passwords hashed with different algorithms can co-exist.

4. Some of them are adaptive — that means when computers get faster, you can tune the algorithm to be slower, without introducing incompatibility with existing passwords.

Table F.16 lists the algorithms supported by the `crypt()` function.
Table F.16. Supported Algorithms for crypt()

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Max Password Length</th>
<th>Adaptive?</th>
<th>Salt Bits</th>
<th>Output Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bf</td>
<td>72</td>
<td>yes</td>
<td>128</td>
<td>60</td>
<td>Blowfish-based, variant 2a</td>
</tr>
<tr>
<td>md5</td>
<td>unlimited</td>
<td>no</td>
<td>48</td>
<td>34</td>
<td>MD5-based crypt</td>
</tr>
<tr>
<td>xdes</td>
<td>8</td>
<td>yes</td>
<td>24</td>
<td>20</td>
<td>Extended DES</td>
</tr>
<tr>
<td>des</td>
<td>8</td>
<td>no</td>
<td>12</td>
<td>13</td>
<td>Original UNIX crypt</td>
</tr>
</tbody>
</table>

F.25.2.1. crypt()

crypt(password text, salt text) returns text

Calculates a crypt(3)-style hash of password. When storing a new password, you need to use gen_salt() to generate a new salt value. To check a password, pass the stored hash value as salt, and test whether the result matches the stored value.

Example of setting a new password:

UPDATE ... SET pswhash = crypt('new password', gen_salt('md5'));

Example of authentication:

SELECT (pswhash = crypt('entered password', pswhash)) AS pswmatch FROM ... ;

This returns true if the entered password is correct.

F.25.2.2. gen_salt()

gen_salt(type text [, iter_count integer ]) returns text

Generates a new random salt string for use in crypt(). The salt string also tells crypt() which algorithm to use.

The type parameter specifies the hashing algorithm. The accepted types are: des, xdes, md5 and bf.

The iter_count parameter lets the user specify the iteration count, for algorithms that have one. The higher the count, the more time it takes to hash the password and therefore the more time to break it. Although with too high a count the time to calculate a hash may be several years — which is somewhat impractical. If the iter_count parameter is omitted, the default iteration count is used. Allowed values for iter_count depend on the algorithm and are shown in Table F.17.

Table F.17. Iteration Counts for crypt()

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Default</th>
<th>Min</th>
<th>Max</th>
</tr>
</thead>
<tbody>
<tr>
<td>xdes</td>
<td>725</td>
<td>1</td>
<td>16777215</td>
</tr>
<tr>
<td>bf</td>
<td>6</td>
<td>4</td>
<td>31</td>
</tr>
</tbody>
</table>
Additional Supplied Modules

For xdes there is an additional limitation that the iteration count must be an odd number.

To pick an appropriate iteration count, consider that the original DES crypt was designed to have the speed of 4 hashes per second on the hardware of that time. Slower than 4 hashes per second would probably dampen usability. Faster than 100 hashes per second is probably too fast.

Table F.18 gives an overview of the relative slowness of different hashing algorithms. The table shows how much time it would take to try all combinations of characters in an 8-character password, assuming that the password contains either only lower case letters, or upper- and lower-case letters and numbers. In the crypt-bf entries, the number after a slash is the iter_count parameter of gen_salt.

Table F.18. Hash Algorithm Speeds

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Hashes/sec</th>
<th>For [a-z]</th>
<th>For [A-Za-z0-9]</th>
<th>Duration relative to md5 hash</th>
</tr>
</thead>
<tbody>
<tr>
<td>crypt-bf/8</td>
<td>1792</td>
<td>4 years</td>
<td>3927 years</td>
<td>100k</td>
</tr>
<tr>
<td>crypt-bf/7</td>
<td>3648</td>
<td>2 years</td>
<td>1929 years</td>
<td>50k</td>
</tr>
<tr>
<td>crypt-bf/6</td>
<td>7168</td>
<td>1 year</td>
<td>982 years</td>
<td>25k</td>
</tr>
<tr>
<td>crypt-bf/5</td>
<td>13504</td>
<td>188 days</td>
<td>521 years</td>
<td>12.5k</td>
</tr>
<tr>
<td>crypt-md5</td>
<td>171584</td>
<td>15 days</td>
<td>41 years</td>
<td>1k</td>
</tr>
<tr>
<td>crypt-des</td>
<td>23221568</td>
<td>157.5 minutes</td>
<td>108 days</td>
<td>7</td>
</tr>
<tr>
<td>sha1</td>
<td>37774272</td>
<td>90 minutes</td>
<td>68 days</td>
<td>4</td>
</tr>
<tr>
<td>md5 (hash)</td>
<td>150085504</td>
<td>22.5 minutes</td>
<td>17 days</td>
<td>1</td>
</tr>
</tbody>
</table>

Notes:

- The machine used is an Intel Mobile Core i3.
- crypt-des and crypt-md5 algorithm numbers are taken from John the Ripper v1.6.38 -test output.
- md5 hash numbers are from mdcrack 1.2.
- sha1 numbers are from lcrack-20031130-beta.
- crypt-bf numbers are taken using a simple program that loops over 1000 8-character passwords. That way I can show the speed with different numbers of iterations. For reference: john -test shows 13506 loops/sec for crypt-bf/5. (The very small difference in results is in accordance with the fact that the crypt-bf implementation in pgcrypto is the same one used in John the Ripper.)

Note that “try all combinations” is not a realistic exercise. Usually password cracking is done with the help of dictionaries, which contain both regular words and various mutations of them. So, even somewhat word-like passwords could be cracked much faster than the above numbers suggest, while a 6-character non-word-like password may escape cracking. Or not.

F.25.3. PGP Encryption Functions

The functions here implement the encryption part of the OpenPGP (RFC 4880) standard. Supported are both symmetric-key and public-key encryption.

An encrypted PGP message consists of 2 parts, or packets:

- Packet containing a session key — either symmetric-key or public-key encrypted.
- Packet containing data encrypted with the session key.

When encrypting with a symmetric key (i.e., a password):
1. The given password is hashed using a String2Key (S2K) algorithm. This is rather similar to crypt() algorithms — purposefully slow and with random salt — but it produces a full-length binary key.

2. If a separate session key is requested, a new random key will be generated. Otherwise the S2K key will be used directly as the session key.

3. If the S2K key is to be used directly, then only S2K settings will be put into the session key packet. Otherwise the session key will be encrypted with the S2K key and put into the session key packet.

When encrypting with a public key:
1. A new random session key is generated.
2. It is encrypted using the public key and put into the session key packet.

In either case the data to be encrypted is processed as follows:
1. Optional data-manipulation: compression, conversion to UTF-8, and/or conversion of line-endings.
2. The data is prefixed with a block of random bytes. This is equivalent to using a random IV.
3. An SHA1 hash of the random prefix and data is appended.
4. All this is encrypted with the session key and placed in the data packet.

**F.25.3.1. pgp_sym_encrypt()**

```
pgp_sym_encrypt(data text, psw text [, options text ]) returns bytea
pgp_sym_encrypt_bytea(data bytea, psw text [, options text ]) returns bytea
```

Encrypt data with a symmetric PGP key psw. The options parameter can contain option settings, as described below.

**F.25.3.2. pgp_sym_decrypt()**

```
pgp_sym_decrypt(msg bytea, psw text [, options text ]) returns text
pgp_sym_decrypt_bytea(msg bytea, psw text [, options text ]) returns bytea
```

Decrypt a symmetric-key-encrypted PGP message.

Decrypting bytea data with pgp_sym_decrypt is disallowed. This is to avoid outputting invalid character data. Decrypting originally textual data with pgp_sym_decrypt_bytea is fine.

The options parameter can contain option settings, as described below.

**F.25.3.3. pgp_pub_encrypt()**

```
p gp_pub_encrypt(data text, key bytea [, options text ]) returns bytea
pgp_pub_encrypt_bytea(data bytea, key bytea [, options text ]) returns bytea
```

Encrypt data with a public PGP key key. Giving this function a secret key will produce an error.

The options parameter can contain option settings, as described below.
F.25.3.4. \textit{pgp\_pub\_decrypt} ()

\begin{verbatim}
pgp_pub_decrypt(msg bytea, key bytea [, psw text [, options text ]]) returns text
pgp_pub_decrypt_bytea(msg bytea, key bytea [, psw text [, options text ]]) returns bytea
\end{verbatim}

Decrypt a public-key-encrypted message. \textit{key} must be the secret key corresponding to the public key that was used to encrypt. If the secret key is password-protected, you must give the password in \textit{psw}. If there is no password, but you want to specify options, you need to give an empty password.

Decrypting \textit{bytea} data with \textit{pgp\_pub\_decrypt} is disallowed. This is to avoid outputting invalid character data. Decrypting originally textual data with \textit{pgp\_pub\_decrypt\_bytea} is fine.

The \textit{options} parameter can contain option settings, as described below.

F.25.3.5. \textit{pgp\_key\_id}()

\begin{verbatim}
pgp_key_id(bytea) returns text
\end{verbatim}

\textit{pgp\_key\_id} extracts the key ID of a PGP public or secret key. Or it gives the key ID that was used for encrypting the data, if given an encrypted message.

It can return 2 special key IDs:

\begin{itemize}
  \item SYMKEY
    The message is encrypted with a symmetric key.
  \item ANYKEY
    The message is public-key encrypted, but the key ID has been removed. That means you will need to try all your secret keys on it to see which one decrypts it. \textit{pgcrypto} itself does not produce such messages.
\end{itemize}

Note that different keys may have the same ID. This is rare but a normal event. The client application should then try to decrypt with each one, to see which fits — like handling ANYKEY.

F.25.3.6. \textit{armor}(), \textit{dearmor}()

\begin{verbatim}
armor(data bytea [, , keys text[], values text[]]) returns text
dearmor(data text) returns bytea
\end{verbatim}

These functions wrap/unwrap binary data into PGP ASCII-armor format, which is basically Base64 with CRC and additional formatting.

If the \textit{keys} and \textit{values} arrays are specified, an armor header is added to the armored format for each key/value pair. Both arrays must be single-dimensional, and they must be of the same length. The keys and values cannot contain any non-ASCII characters.

F.25.3.7. \textit{pgp\_armor\_headers}

\begin{verbatim}
pgp_armor_headers(data text, key out text, value out text) returns setof record
pgp_armor_headers() extracts the armor headers from \textit{data}. The return value is a set of rows with two columns, key and value. If the keys or values contain any non-ASCII characters, they are treated as UTF-8.
\end{verbatim}

3388
F.25.3.8. Options for PGP Functions

Options are named to be similar to GnuPG. An option's value should be given after an equal sign; separate options from each other with commas. For example:

```plaintext
pgp_sym_encrypt(data, psw, 'compress-algo=1, cipher-algo=aes256')
```

All of the options except `convert-crlf` apply only to encrypt functions. Decrypt functions get the parameters from the PGP data.

The most interesting options are probably `compress-algo` and `unicode-mode`. The rest should have reasonable defaults.

F.25.3.8.1. cipher-algo

Which cipher algorithm to use.

Values: bf, aes128, aes192, aes256 (OpenSSL-only: 3des, cast5)
Default: aes128
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

F.25.3.8.2. compress-algo

Which compression algorithm to use. Only available if PostgreSQL was built with zlib.

Values:
- 0 - no compression
- 1 - ZIP compression
- 2 - ZLIB compression (= ZIP plus meta-data and block CRCs)
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

F.25.3.8.3. compress-level

How much to compress. Higher levels compress smaller but are slower. 0 disables compression.

Values: 0, 1-9
Default: 6
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

F.25.3.8.4. convert-crlf

Whether to convert \n into \r\n when encrypting and \r\n to \n when decrypting. RFC 4880 specifies that text data should be stored using \r\n line-feeds. Use this to get fully RFC-compliant behavior.

Values: 0, 1
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt, pgp_sym_decrypt, pgp_pub_decrypt

F.25.3.8.5. disable-mdc

Do not protect data with SHA-1. The only good reason to use this option is to achieve compatibility with ancient PGP products, predating the addition of SHA-1 protected packets to RFC 4880. Recent gnupg.org and pgp.com software supports it fine.


**Values: 0, 1**  
**Default: 0**  
**Applies to: pgp_sym_encrypt, pgp_pub_encrypt**

**F.25.3.8.6. sess-key**

Use separate session key. Public-key encryption always uses a separate session key; this option is for  
symmetric-key encryption, which by default uses the S2K key directly.

**Values: 0, 1**  
**Default: 0**  
**Applies to: pgp_sym_encrypt**

**F.25.3.8.7. s2k-mode**

Which S2K algorithm to use.

**Values:**  
0 - Without salt. Dangerous!  
1 - With salt but with fixed iteration count.  
3 - Variable iteration count.  
**Default: 3**  
**Applies to: pgp_sym_encrypt**

**F.25.3.8.8. s2k-count**

The number of iterations of the S2K algorithm to use. It must be a value between 1024 and 65011712,  
inclusive.

**Default: A random value between 65536 and 253952**  
**Applies to: pgp_sym_encrypt, only with s2k-mode=3**

**F.25.3.8.9. s2k-digest-algo**

Which digest algorithm to use in S2K calculation.

**Values: md5, sha1**  
**Default: sha1**  
**Applies to: pgp_sym_encrypt**

**F.25.3.8.10. s2k-cipher-algo**

Which cipher to use for encrypting separate session key.

**Values: bf, aes, aes128, aes192, aes256**  
**Default: use cipher-algo**  
**Applies to: pgp_sym_encrypt**

**F.25.3.8.11. unicode-mode**

Whether to convert textual data from database internal encoding to UTF-8 and back. If your database  
already is UTF-8, no conversion will be done, but the message will be tagged as UTF-8. Without this  
option it will not be.
F.25.3.9. Generating PGP Keys with GnuPG

To generate a new key:

```
gpg --gen-key
```

The preferred key type is “DSA and Elgamal”.

For RSA encryption you must create either DSA or RSA sign-only key as master and then add an
RSA encryption subkey with `gpg --edit-key`.

To list keys:

```
gpg --list-secret-keys
```

To export a public key in ASCII-armor format:

```
gpg -a --export KEYID > public.key
```

To export a secret key in ASCII-armor format:

```
gpg -a --export-secret-keys KEYID > secret.key
```

You need to use `dearmor()` on these keys before giving them to the PGP functions. Or if you can
handle binary data, you can drop `-a` from the command.

For more details see `man gpg`, The GNU Privacy Handbook[^3] and other documentation on https://
www.gnupg.org/.

F.25.3.10. Limitations of PGP Code

- No support for signing. That also means that it is not checked whether the encryption subkey belongs
to the master key.

- No support for encryption key as master key. As such practice is generally discouraged, this should
not be a problem.

- No support for several subkeys. This may seem like a problem, as this is common practice. On the
other hand, you should not use your regular GPG/PGP keys with `pgcrypto`, but create new ones,
as the usage scenario is rather different.

F.25.4. Raw Encryption Functions

These functions only run a cipher over data; they don't have any advanced features of PGP encryption.
Therefore they have some major problems:

1. They use user key directly as cipher key.

2. They don't provide any integrity checking, to see if the encrypted data was modified.

3. They expect that users manage all encryption parameters themselves, even IV.

4. They don’t handle text.

So, with the introduction of PGP encryption, usage of raw encryption functions is discouraged.

```sql
encrypt(data bytea, key bytea, type text) returns bytea
decrypt(data bytea, key bytea, type text) returns bytea
encrypt_iv(data bytea, key bytea, iv bytea, type text) returns bytea
decrypt_iv(data bytea, key bytea, iv bytea, type text) returns bytea
```

Encrypt/decrypt data using the cipher method specified by `type`. The syntax of the `type` string is:

```
algorithm [ - mode ] [ /pad: padding ]
```

where `algorithm` is one of:

- `bf` — Blowfish
- `aes` — AES (Rijndael-128, -192 or -256)

and `mode` is one of:

- `cbc` — next block depends on previous (default)
- `ecb` — each block is encrypted separately (for testing only)

and `padding` is one of:

- `pkcs` — data may be any length (default)
- `none` — data must be multiple of cipher block size

So, for example, these are equivalent:

```sql
encrypt(data, 'fooz', 'bf')
encrypt(data, 'fooz', 'bf-cbc/pad:pkcs')
```

In `encrypt_iv` and `decrypt_iv`, the `iv` parameter is the initial value for the CBC mode; it is ignored for ECB. It is clipped or padded with zeroes if not exactly block size. It defaults to all zeroes in the functions without this parameter.

**F.25.5. Random-Data Functions**

```sql
gen_random_bytes(count integer) returns bytea
```

Returns `count` cryptographically strong random bytes. At most 1024 bytes can be extracted at a time. This is to avoid draining the randomness generator pool.

```sql
gen_random_uuid() returns uuid
```

Returns a version 4 (random) UUID.
**F.25.6. Notes**

**F.25.6.1. Configuration**

`pgcrypto` configures itself according to the findings of the main PostgreSQL `configure` script. The options that affect it are `--with-zlib` and `--with-openssl`.

When compiled with zlib, PGP encryption functions are able to compress data before encrypting.

When compiled with OpenSSL, there will be more algorithms available. Also public-key encryption functions will be faster as OpenSSL has more optimized BIGNUM functions.

<table>
<thead>
<tr>
<th>Functionality</th>
<th>Built-in</th>
<th>With OpenSSL</th>
</tr>
</thead>
<tbody>
<tr>
<td>MD5</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>SHA1</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>SHA224/256/384/512</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>Other digest algorithms</td>
<td>no</td>
<td>yes (Note 1)</td>
</tr>
<tr>
<td>Blowfish</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>AES</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>DES/3DES/CAST5</td>
<td>no</td>
<td>yes</td>
</tr>
<tr>
<td>Raw encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>PGP Symmetric encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>PGP Public-Key encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
</tbody>
</table>

Notes:

1. Any digest algorithm OpenSSL supports is automatically picked up. This is not possible with ciphers, which need to be supported explicitly.

**F.25.6.2. NULL Handling**

As is standard in SQL, all functions return NULL, if any of the arguments are NULL. This may create security risks on careless usage.

**F.25.6.3. Security Limitations**

All `pgcrypto` functions run inside the database server. That means that all the data and passwords move between `pgcrypto` and client applications in clear text. Thus you must:

1. Connect locally or use SSL connections.
2. Trust both system and database administrator.

If you cannot, then better do crypto inside client application.

The implementation does not resist side-channel attacks\(^4\). For example, the time required for a `pgcrypto` decryption function to complete varies among ciphertexts of a given size.

**F.25.6.4. Useful Reading**


  The GNU Privacy Handbook.

---

\(^4\) [https://en.wikipedia.org/wiki/Side-channel_attack](https://en.wikipedia.org/wiki/Side-channel_attack)
Additional Supplied Modules

  Describes the crypt-blowfish algorithm.
  How to choose a good password.
  Interesting idea for picking passwords.
- [http://www.interhack.net/people/cmcurtin/snake-oil-faq.html](http://www.interhack.net/people/cmcurtin/snake-oil-faq.html)
  Describes good and bad cryptography.

### F.25.6.5. Technical References

  OpenPGP message format.
  The MD5 Message-Digest Algorithm.
  HMAC: Keyed-Hashing for Message Authentication.
- [https://www.usenix.org/legacy/events/usenix99/provos.html](https://www.usenix.org/legacy/events/usenix99/provos.html)
  Comparison of crypt-des, crypt-md5 and bcrypt algorithms.
  Description of Fortuna CSPRNG.
- [http://jlcooke.ca/random/](http://jlcooke.ca/random/)
  Jean-Luc Cooke Fortuna-based /dev/random driver for Linux.

### F.25.7. Author

Marko Kreen <markokr@gmail.com>

pgcrypto uses code from the following sources:

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Author</th>
<th>Source origin</th>
</tr>
</thead>
<tbody>
<tr>
<td>DES crypt</td>
<td>David Burren and others</td>
<td>FreeBSD libcrypt</td>
</tr>
<tr>
<td>MD5 crypt</td>
<td>Poul-Henning Kamp</td>
<td>FreeBSD libcrypt</td>
</tr>
<tr>
<td>Blowfish crypt</td>
<td>Solar Designer</td>
<td><a href="http://www.openwall.com">www.openwall.com</a></td>
</tr>
<tr>
<td>Blowfish cipher</td>
<td>Simon Tatham</td>
<td>PuTTY</td>
</tr>
<tr>
<td>Rijndael cipher</td>
<td>Brian Gladman</td>
<td>OpenBSD sys/crypto</td>
</tr>
<tr>
<td>MD5 hash and SHA1</td>
<td>WIDE Project</td>
<td>KAME kame/sys/crypto</td>
</tr>
<tr>
<td>SHA256/384/512</td>
<td>Aaron D. Gifford</td>
<td>OpenBSD sys/crypto</td>
</tr>
<tr>
<td>BIGNUM math</td>
<td>Michael J. Fromberger</td>
<td>dartmouth.edu/~sting/sw/imath</td>
</tr>
</tbody>
</table>
F.26. pg_freespacemap

The pg_freespacemap module provides a means for examining the free space map (FSM). It provides a function called pg_freespace, or two overloaded functions, to be precise. The functions show the value recorded in the free space map for a given page, or for all pages in the relation.

By default use is restricted to superusers and members of the pg_stat_scan_tables role. Access may be granted to others using GRANT.

F.26.1. Functions

\[
\text{pg_freespace}\text{(rel regclass IN, blkno bigint IN) returns int2}
\]

Returns the amount of free space on the page of the relation, specified by blkno, according to the FSM.

\[
\text{pg_freespace}\text{(rel regclass IN, blkno OUT bigint, avail OUT int2)}
\]

Displays the amount of free space on each page of the relation, according to the FSM. A set of (blkno bigint, avail int2) tuples is returned, one tuple for each page in the relation.

The values stored in the free space map are not exact. They're rounded to precision of $1/256$th of BLCKSZ (32 bytes with default BLCKSZ), and they're not kept fully up-to-date as tuples are inserted and updated.

For indexes, what is tracked is entirely-unused pages, rather than free space within pages. Therefore, the values are not meaningful, just whether a page is full or empty.

Note

The interface was changed in version 8.4, to reflect the new FSM implementation introduced in the same version.

F.26.2. Sample Output

\[
\text{postgres=# SELECT * FROM pg_freespace('foo');}
\]

<table>
<thead>
<tr>
<th>blkno</th>
<th>avail</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>2</td>
<td>0</td>
</tr>
<tr>
<td>3</td>
<td>32</td>
</tr>
<tr>
<td>4</td>
<td>704</td>
</tr>
<tr>
<td>5</td>
<td>704</td>
</tr>
<tr>
<td>6</td>
<td>704</td>
</tr>
<tr>
<td>7</td>
<td>1216</td>
</tr>
<tr>
<td>8</td>
<td>704</td>
</tr>
<tr>
<td>9</td>
<td>704</td>
</tr>
<tr>
<td>10</td>
<td>704</td>
</tr>
<tr>
<td>11</td>
<td>704</td>
</tr>
<tr>
<td>12</td>
<td>704</td>
</tr>
<tr>
<td>13</td>
<td>704</td>
</tr>
<tr>
<td>14</td>
<td>704</td>
</tr>
<tr>
<td>15</td>
<td>704</td>
</tr>
<tr>
<td>16</td>
<td>704</td>
</tr>
</tbody>
</table>
F.26.3. Author

Original version by Mark Kirkwood <markir@paradise.net.nz>. Rewritten in version 8.4 to suit new FSM implementation by Heikki Linnakangas <heikki@enterprisedb.com>

F.27. pg_prewarm

The pg_prewarm module provides a convenient way to load relation data into either the operating system buffer cache or the PostgreSQL buffer cache. Prewarming can be performed manually using the pg_prewarm function, or can be performed automatically by including pg_prewarm in shared_preload_libraries. In the latter case, the system will run a background worker which periodically records the contents of shared buffers in a file called autoprewarm.blocks and will, using 2 background workers, reload those same blocks after a restart.

F.27.1. Functions

```
pg_prewarm(regclass, mode text default 'buffer', fork text default 'main',
            first_block int8 default null,
            last_block int8 default null) RETURNS int8
```

The first argument is the relation to be prewarmed. The second argument is the prewarming method to be used, as further discussed below; the third is the relation fork to be prewarmed, usually main. The fourth argument is the first block number to prewarm (NULL is accepted as a synonym for zero). The fifth argument is the last block number to prewarm (NULL means prewarm through the last block in the relation). The return value is the number of blocks prewarmed.

There are three available prewarming methods. prefetch issues asynchronous prefetch requests to the operating system, if this is supported, or throws an error otherwise. read reads the requested range of blocks; unlike prefetch, this is synchronous and supported on all platforms and builds, but may be slower. buffer reads the requested range of blocks into the database buffer cache.

Note that with any of these methods, attempting to prewarm more blocks than can be cached — by the OS when using prefetch or read, or by PostgreSQL when using buffer — will likely result in lower-numbered blocks being evicted as higher numbered blocks are read in. Prewarmed data also enjoys no special protection from cache evictions, so it is possible that other system activity may evict the newly prewarmed blocks shortly after they are read; conversely, prewarming may also evict other data from cache. For these reasons, prewarming is typically most useful at startup, when caches are largely empty.

```
autoprewarm_start_worker() RETURNS void
```

Launch the main autoprewarm worker. This will normally happen automatically, but is useful if automatic prewarm was not configured at server startup time and you wish to start up the worker at a later time.
autoprewarm_dump_now() RETURNS int8

Update autoprewarm.blocks immediately. This may be useful if the autoprewarm worker is not running but you anticipate running it after the next restart. The return value is the number of records written to autoprewarm.blocks.

F.27.2. Configuration Parameters

pg_prewarm.autoprewarm(boolean)

Controls whether the server should run the autoprewarm worker. This is on by default. This parameter can only be set at server start.

pg_prewarm.autoprewarm_interval(int)

This is the interval between updates to autoprewarm.blocks. The default is 300 seconds. If set to 0, the file will not be dumped at regular intervals, but only when the server is shut down.

F.27.3. Author

Robert Haas <rhaas@postgresql.org>

F.28. pgrowlocks

The pgrowlocks module provides a function to show row locking information for a specified table.

By default use is restricted to superusers, members of the pg_stat_scan_tables role, and users with SELECT permissions on the table.

F.28.1. Overview

pgrowlocks(text) returns setof record

The parameter is the name of a table. The result is a set of records, with one row for each locked row within the table. The output columns are shown in Table F.20.

Table F.20. pgrowlocks Output Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>locked_row</td>
<td>tid</td>
<td>Tuple ID (TID) of locked row</td>
</tr>
<tr>
<td>locker</td>
<td>xid</td>
<td>Transaction ID of locker, or multixact ID if multitransaction</td>
</tr>
<tr>
<td>multi</td>
<td>boolean</td>
<td>True if locker is a multitransaction</td>
</tr>
<tr>
<td>xids</td>
<td>xid[]</td>
<td>Transaction IDs of lockers (more than one if multitransaction)</td>
</tr>
<tr>
<td>modes</td>
<td>text[]</td>
<td>Lock mode of lockers (more than one if multitransaction), an array of Key Share, Share, For No Key Update, No Key Update, For Update, Update.</td>
</tr>
</tbody>
</table>
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pids</td>
<td>integer[]</td>
<td>Process IDs of locking backends (more than one if multitransaction)</td>
</tr>
</tbody>
</table>

`pgrowlocks` takes `AccessShareLock` for the target table and reads each row one by one to collect the row locking information. This is not very speedy for a large table. Note that:

1. If the table as a whole is exclusive-locked by someone else, `pgrowlocks` will be blocked.
2. `pgrowlocks` is not guaranteed to produce a self-consistent snapshot. It is possible that a new row lock is taken, or an old lock is freed, during its execution.

`pgrowlocks` does not show the contents of locked rows. If you want to take a look at the row contents at the same time, you could do something like this:

```sql
SELECT * FROM accounts AS a, pgrowlocks('accounts') AS p
WHERE p.locked_row = a.ctid;
```

Be aware however that such a query will be very inefficient.

#### F.28.2. Sample Output

```sql
=# SELECT * FROM pgrowlocks('t1');

<table>
<thead>
<tr>
<th>locked_row</th>
<th>locker</th>
<th>multi</th>
<th>xids</th>
<th>modes</th>
<th>pids</th>
</tr>
</thead>
<tbody>
<tr>
<td>(0,1)</td>
<td>609</td>
<td>f</td>
<td>{609}</td>
<td>&quot;For Share&quot;</td>
<td>{3161}</td>
</tr>
<tr>
<td>(0,2)</td>
<td>609</td>
<td>f</td>
<td>{609}</td>
<td>&quot;For Share&quot;</td>
<td>{3161}</td>
</tr>
<tr>
<td>(0,3)</td>
<td>607</td>
<td>f</td>
<td>{607}</td>
<td>&quot;For Update&quot;</td>
<td>{3107}</td>
</tr>
<tr>
<td>(0,4)</td>
<td>607</td>
<td>f</td>
<td>{607}</td>
<td>&quot;For Update&quot;</td>
<td>{3107}</td>
</tr>
</tbody>
</table>
```

(4 rows)

#### F.28.3. Author

Tatsuo Ishii

### F.29. `pg_stat_statements`

The `pg_stat_statements` module provides a means for tracking execution statistics of all SQL statements executed by a server.

The module must be loaded by adding `pg_stat_statements` to `shared_preload_libraries` in `postgresql.conf`, because it requires additional shared memory. This means that a server restart is needed to add or remove the module.

When `pg_stat_statements` is loaded, it tracks statistics across all databases of the server. To access and manipulate these statistics, the module provides a view, `pg_stat_statements`, and the utility functions `pg_stat_statements_reset` and `pg_stat_statements`. These are not available globally but can be enabled for a specific database with `CREATE EXTENSION pg_stat_statements`.

#### F.29.1. The `pg_stat_statements` View

The statistics gathered by the module are made available via a view named `pg_stat_statements`. This view contains one row for each distinct database ID, user ID and query ID (up to the maximum
number of distinct statements that the module can track). The columns of the view are shown in Table F.21.

**Table F.21. pg_stat_statements Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>userid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>OID of user who executed the statement</td>
</tr>
<tr>
<td>dbid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>OID of database in which the statement was executed</td>
</tr>
<tr>
<td>queryid</td>
<td>bigint</td>
<td></td>
<td>Internal hash code, computed from the statement's parse tree</td>
</tr>
<tr>
<td>query</td>
<td>text</td>
<td></td>
<td>Text of a representative statement</td>
</tr>
<tr>
<td>calls</td>
<td>bigint</td>
<td></td>
<td>Number of times executed</td>
</tr>
<tr>
<td>total_time</td>
<td>double precision</td>
<td></td>
<td>Total time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>min_time</td>
<td>double precision</td>
<td></td>
<td>Minimum time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>max_time</td>
<td>double precision</td>
<td></td>
<td>Maximum time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>mean_time</td>
<td>double precision</td>
<td></td>
<td>Mean time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>stddev_time</td>
<td>double precision</td>
<td></td>
<td>Population standard deviation of time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>rows</td>
<td>bigint</td>
<td></td>
<td>Total number of rows retrieved or affected by the statement</td>
</tr>
<tr>
<td>shared_blks_hit</td>
<td>bigint</td>
<td></td>
<td>Total number of shared block cache hits by the statement</td>
</tr>
<tr>
<td>shared_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks read by the statement</td>
</tr>
<tr>
<td>shared_blks_dirtied</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks dirtied by the statement</td>
</tr>
<tr>
<td>shared_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks written by the statement</td>
</tr>
<tr>
<td>local_blks_hit</td>
<td>bigint</td>
<td></td>
<td>Total number of local block cache hits by the statement</td>
</tr>
</tbody>
</table>

3399
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>local_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks read by the statement</td>
</tr>
<tr>
<td>local_blks_dirtied</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks dirtied by the statement</td>
</tr>
<tr>
<td>local_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks written by the statement</td>
</tr>
<tr>
<td>temp_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of temp blocks read by the statement</td>
</tr>
<tr>
<td>temp_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of temp blocks written by the statement</td>
</tr>
<tr>
<td>blk_read_time</td>
<td>double precision</td>
<td></td>
<td>Total time the statement spent reading blocks, in milliseconds (if track_io_timing is enabled, otherwise zero)</td>
</tr>
<tr>
<td>blk_write_time</td>
<td>double precision</td>
<td></td>
<td>Total time the statement spent writing blocks, in milliseconds (if track_io_timing is enabled, otherwise zero)</td>
</tr>
</tbody>
</table>

For security reasons, only superusers and members of the `pg_read_all_stats` role are allowed to see the SQL text and `queryid` of queries executed by other users. Other users can see the statistics, however, if the view has been installed in their database.

Plannable queries (that is, `SELECT`, `INSERT`, `UPDATE`, and `DELETE`) are combined into a single `pg_stat_statements` entry whenever they have identical query structures according to an internal hash calculation. Typically, two queries will be considered the same for this purpose if they are semantically equivalent except for the values of literal constants appearing in the query. Utility commands (that is, all other commands) are compared strictly on the basis of their textual query strings, however.

When a constant’s value has been ignored for purposes of matching the query to other queries, the constant is replaced by a parameter symbol, such as $1$, in the `pg_stat_statements` display. The rest of the query text is that of the first query that had the particular `queryid` hash value associated with the `pg_stat_statements` entry.

In some cases, queries with visibly different texts might get merged into a single `pg_stat_statements` entry. Normally this will happen only for semantically equivalent queries, but there is a small chance of hash collisions causing unrelated queries to be merged into one entry. (This cannot happen for queries belonging to different users or databases, however.)

Since the `queryid` hash value is computed on the post-parse-analysis representation of the queries, the opposite is also possible: queries with identical texts might appear as separate entries, if they have different meanings as a result of factors such as different `search_path` settings.

Consumers of `pg_stat_statements` may wish to use `queryid` (perhaps in combination with `dbid` and `userid`) as a more stable and reliable identifier for each entry than its query text. How-
ever, it is important to understand that there are only limited guarantees around the stability of the queryid hash value. Since the identifier is derived from the post-parse-analysis tree, its value is a function of, among other things, the internal object identifiers appearing in this representation. This has some counterintuitive implications. For example, pg_stat_statements will consider two apparently-identical queries to be distinct, if they reference a table that was dropped and recreated between the executions of the two queries. The hashing process is also sensitive to differences in machine architecture and other facets of the platform. Furthermore, it is not safe to assume that queryid will be stable across major versions of PostgreSQL.

As a rule of thumb, queryid values can be assumed to be stable and comparable only so long as the underlying server version and catalog metadata details stay exactly the same. Two servers participating in replication based on physical WAL replay can be expected to have identical queryid values for the same query. However, logical replication schemes do not promise to keep replicas identical in all relevant details, so queryid will not be a useful identifier for accumulating costs across a set of logical replicas. If in doubt, direct testing is recommended.

The parameter symbols used to replace constants in representative query texts start from the next number after the highest $n$ parameter in the original query text, or $1$ if there was none. It's worth noting that in some cases there may be hidden parameter symbols that affect this numbering. For example, PL/pgSQL uses hidden parameter symbols to insert values of function local variables into queries, so that a PL/pgSQL statement like SELECT i + 1 INTO j would have representative text like SELECT i + $2$.

The representative query texts are kept in an external disk file, and do not consume shared memory. Therefore, even very lengthy query texts can be stored successfully. However, if many long query texts are accumulated, the external file might grow unmanageably large. As a recovery method if that happens, pg_stat_statements may choose to discard the query texts, whereupon all existing entries in the pg_stat_statements view will show null query fields, though the statistics associated with each queryid are preserved. If this happens, consider reducing pg_stat_statements.max to prevent recurrences.

**F.29.2. Functions**

pg_stat_statements_reset() returns void

pg_stat_statements_reset discards all statistics gathered so far by pg_stat_statements. By default, this function can only be executed by superusers.

pg_stat_statements(showtext boolean) returns setof record

The pg_stat_statements view is defined in terms of a function also named pg_stat_statements. It is possible for clients to call the pg_stat_statements function directly, and by specifying showtext := false have query text be omitted (that is, the OUT argument that corresponds to the view's query column will return nulls). This feature is intended to support external tools that might wish to avoid the overhead of repeatedly retrieving query texts of indeterminate length. Such tools can instead cache the first query text observed for each entry themselves, since that is all pg_stat_statements itself does, and then retrieve query texts only as needed. Since the server stores query texts in a file, this approach may reduce physical I/O for repeated examination of the pg_stat_statements data.

**F.29.3. Configuration Parameters**

pg_stat_statements.max(integer)

pg_stat_statements.max is the maximum number of statements tracked by the module (i.e., the maximum number of rows in the pg_stat_statements view). If more distinct statements than that are observed, information about the least-executed statements is discarded. The default value is 5000. This parameter can only be set at server start.
pg_stat_statements.track (enum)

pg_stat_statements.track controls which statements are counted by the module. Specify top to track top-level statements (those issued directly by clients), all to also track nested statements (such as statements invoked within functions), or none to disable statement statistics collection. The default value is top. Only superusers can change this setting.

pg_stat_statements.track_utility (boolean)

pg_stat_statements.track_utility controls whether utility commands are tracked by the module. Utility commands are all those other than SELECT, INSERT, UPDATE and DELETE. The default value is on. Only superusers can change this setting.

pg_stat_statements.save (boolean)

pg_stat_statements.save specifies whether to save statement statistics across server shutdowns. If it is off then statistics are not saved at shutdown nor reloaded at server start. The default value is on. This parameter can only be set in the postgresql.conf file or on the server command line.

The module requires additional shared memory proportional to pg_stat_statements.max. Note that this memory is consumed whenever the module is loaded, even if pg_stat_statements.track is set to none.

These parameters must be set in postgresql.conf. Typical usage might be:

```
# postgresql.conf
shared_preload_libraries = 'pg_stat_statements'
pg_stat_statements.max = 10000
pg_stat_statements.track = all
```

F.29.4. Sample Output

```
bench=# SELECT pg_stat_statements_reset();

$ pgbench -i bench
$ pgbench -c10 -t300 bench

bench=# \x

bench=# SELECT query, calls, total_time, rows, 100.0 *
shared_blks_hit / nullif(shared_blks_hit + shared_blks_read, 0) AS
hit_percent
FROM pg_stat_statements ORDER BY total_time DESC LIMIT 5;

-+----------------------------------------------------------+
1 | UPDATE pgbench_branches SET bbalance = bbalance + $1    |
 | WHERE bid = $2;                                         |
 | calls | 3000                                                   |
 | total_time | 9609.00100000002                                       |
 | rows | 2836                                                   |
 | hit_percent | 99.9778970000200936                                     |

-+----------------------------------------------------------+
2 | UPDATE pgbench_tellers SET tbalance = tbalance + $1      |
 | WHERE tid = $2;                                         |
```
F.30. pgstattuple

The `pgstattuple` module provides various functions to obtain tuple-level statistics.

Because these functions return detailed page-level information, access is restricted by default. By default, only the role `pg_stat_scan_tables` has EXECUTE privilege. Superusers of course bypass this restriction. After the extension has been installed, users may issue GRANT commands to change the privileges on the functions to allow others to execute them. However, it might be preferable to add those users to the `pg_stat_scan_tables` role instead.

F.30.1. Functions

`pgstattuple(regclass) returns record`

`pgstattuple` returns a relation's physical length, percentage of “dead” tuples, and other info. This may help users to determine whether vacuum is necessary or not. The argument is the target relation's name (optionally schema-qualified) or OID. For example:

```sql
test=> SELECT * FROM pgstattuple('pg_catalog.pg_proc');
- [ RECORD 1 ]---------
  table_len | 458752
  tuple_count | 1470
  tuple_len | 438896
  tuple_percent | 95.67
```

F.29.5. Authors

Takahiro Itagaki <itagaki.takahiro@oss.ntt.co.jp>. Query normalization added by Peter Geoghegan <peter@2ndquadrant.com>.
dead_tuple_count | 11
dead_tuple_len   | 3157
dead_tuple_percent | 0.69
free_space       | 8932
free_percent     | 1.95

The output columns are described in Table F.22.

**Table F.22. pgstattuple Output Columns**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_len</td>
<td>bigint</td>
<td>Physical relation length in bytes</td>
</tr>
<tr>
<td>tuple_count</td>
<td>bigint</td>
<td>Number of live tuples</td>
</tr>
<tr>
<td>tuple_len</td>
<td>bigint</td>
<td>Total length of live tuples in bytes</td>
</tr>
<tr>
<td>tuple_percent</td>
<td>float8</td>
<td>Percentage of live tuples</td>
</tr>
<tr>
<td>dead_tuple_count</td>
<td>bigint</td>
<td>Number of dead tuples</td>
</tr>
<tr>
<td>dead_tuple_len</td>
<td>bigint</td>
<td>Total length of dead tuples in bytes</td>
</tr>
<tr>
<td>dead_tuple_percent</td>
<td>float8</td>
<td>Percentage of dead tuples</td>
</tr>
<tr>
<td>free_space</td>
<td>bigint</td>
<td>Total free space in bytes</td>
</tr>
<tr>
<td>free_percent</td>
<td>float8</td>
<td>Percentage of free space</td>
</tr>
</tbody>
</table>

**Note**

The `table_len` will always be greater than the sum of the `tuple_len`, `dead_tuple_len` and `free_space`. The difference is accounted for by fixed page overhead, the per-page table of pointers to tuples, and padding to ensure that tuples are correctly aligned.

`pgstattuple` acquires only a read lock on the relation. So the results do not reflect an instantaneous snapshot; concurrent updates will affect them.

`pgstattuple` judges a tuple is “dead” if `HeapTupleSatisfiesDirty` returns false.

`pgstattuple(text)` returns record

This is the same as `pgstattuple(regclass)`, except that the target relation is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

`pgstatindex(regclass)` returns record

`pgstatindex` returns a record showing information about a B-tree index. For example:

```
test=> SELECT * FROM pgstatindex('pg_cast_oid_index');
- RECORD 1 |---------
version    | 2
 tree_level | 0
 index_size | 16384
 root_block_no | 1
```
The output columns are:

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>version</td>
<td>integer</td>
<td>B-tree version number</td>
</tr>
<tr>
<td>tree_level</td>
<td>integer</td>
<td>Tree level of the root page</td>
</tr>
<tr>
<td>index_size</td>
<td>bigint</td>
<td>Total index size in bytes</td>
</tr>
<tr>
<td>root_block_no</td>
<td>bigint</td>
<td>Location of root page (zero if none)</td>
</tr>
<tr>
<td>internal_pages</td>
<td>bigint</td>
<td>Number of “internal” (upper-level) pages</td>
</tr>
<tr>
<td>leaf_pages</td>
<td>bigint</td>
<td>Number of leaf pages</td>
</tr>
<tr>
<td>empty_pages</td>
<td>bigint</td>
<td>Number of empty pages</td>
</tr>
<tr>
<td>deleted_pages</td>
<td>bigint</td>
<td>Number of deleted pages</td>
</tr>
<tr>
<td>avg_leaf_density</td>
<td>float8</td>
<td>Average density of leaf pages</td>
</tr>
<tr>
<td>leaf_fragmentation</td>
<td>float8</td>
<td>Leaf page density</td>
</tr>
</tbody>
</table>

The reported `index_size` will normally correspond to one more page than is accounted for by `internal_pages + leaf_pages + empty_pages + deleted_pages`, because it also includes the index's metapage.

As with `pgstattuple`, the results are accumulated page-by-page, and should not be expected to represent an instantaneous snapshot of the whole index.

`pgstatindex(text)` returns record

This is the same as `pgstatindex(regclass)`, except that the target index is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

`pgstatginindex(regclass)` returns record

`pgstatginindex` returns a record showing information about a GIN index. For example:

```
test=> SELECT * FROM pgstatginindex('test_gin_index');
-[ RECORD 1 ]--
version | 1
pending_pages | 0
pending_tuples | 0
```

The output columns are:

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>version</td>
<td>integer</td>
<td>GIN version number</td>
</tr>
<tr>
<td>pending_pages</td>
<td>integer</td>
<td>Number of pages in the pending list</td>
</tr>
</tbody>
</table>
pgstattashashindex(regclass) returns record

pgstattashashindex returns a record showing information about a HASH index. For example:

test=> select * from pgstattashashindex('con_hash_index');
- [ RECORD 1 ]---------------------
   version | 4
   bucket_pages | 33081
   overflow_pages | 0
   bitmap_pages | 1
   unused_pages | 32455
   live_items | 10204006
   dead_items | 0
   free_percent | 61.8005949100872

The output columns are:

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>version</td>
<td>integer</td>
<td>HASH version number</td>
</tr>
<tr>
<td>bucket_pages</td>
<td>bigint</td>
<td>Number of bucket pages</td>
</tr>
<tr>
<td>overflow_pages</td>
<td>bigint</td>
<td>Number of overflow pages</td>
</tr>
<tr>
<td>bitmap_pages</td>
<td>bigint</td>
<td>Number of bitmap pages</td>
</tr>
<tr>
<td>unused_pages</td>
<td>bigint</td>
<td>Number of unused pages</td>
</tr>
<tr>
<td>live_items</td>
<td>bigint</td>
<td>Number of live tuples</td>
</tr>
<tr>
<td>dead_items</td>
<td>bigint</td>
<td>Number of dead tuples</td>
</tr>
<tr>
<td>free_percent</td>
<td>float</td>
<td>Percentage of free space</td>
</tr>
</tbody>
</table>

pg_relpages(regclass) returns bigint

pg_relpages returns the number of pages in the relation.

pg_relpages(text) returns bigint

This is the same as pg_relpages(regclass), except that the target relation is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

pgstattuple_approx(regclass) returns record

pgstattuple_approx is a faster alternative to pgstattuple that returns approximate results. The argument is the target relation's name or OID. For example:

test=> SELECT * FROM pgstattuple_approx('pg_catalog.pg_proc'::regclass);
- [ RECORD 1 ]---------------------
   table_len | 573440
   scanned_percent | 2
   approx_tuple_count | 2740
   approx_tuple_len | 561210
   approx_tuple_percent | 97.87
The output columns are described in Table F.23.

Whereas pgstattuple always performs a full-table scan and returns an exact count of live and dead tuples (and their sizes) and free space, pgstattuple_approx tries to avoid the full-table scan and returns exact dead tuple statistics along with an approximation of the number and size of live tuples and free space.

It does this by skipping pages that have only visible tuples according to the visibility map (if a page has the corresponding VM bit set, then it is assumed to contain no dead tuples). For such pages, it derives the free space value from the free space map, and assumes that the rest of the space on the page is taken up by live tuples.

For pages that cannot be skipped, it scans each tuple, recording its presence and size in the appropriate counters, and adding up the free space on the page. At the end, it estimates the total number of live tuples based on the number of pages and tuples scanned (in the same way that VACUUM estimates pg_class.reltuples).

**Table F.23. pgstattuple_approx Output Columns**

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_len</td>
<td>bigint</td>
<td>Physical relation length in bytes (exact)</td>
</tr>
<tr>
<td>scanned_percent</td>
<td>float8</td>
<td>Percentage of table scanned</td>
</tr>
<tr>
<td>approx_tuple_count</td>
<td>bigint</td>
<td>Number of live tuples (estimated)</td>
</tr>
<tr>
<td>approx_tuple_len</td>
<td>bigint</td>
<td>Total length of live tuples in bytes (estimated)</td>
</tr>
<tr>
<td>approx_tuple_percent</td>
<td>float8</td>
<td>Percentage of live tuples</td>
</tr>
<tr>
<td>dead_tuple_count</td>
<td>bigint</td>
<td>Number of dead tuples (exact)</td>
</tr>
<tr>
<td>dead_tuple_len</td>
<td>bigint</td>
<td>Total length of dead tuples in bytes (exact)</td>
</tr>
<tr>
<td>dead_tuple_percent</td>
<td>float8</td>
<td>Percentage of dead tuples</td>
</tr>
<tr>
<td>approx_free_space</td>
<td>bigint</td>
<td>Total free space in bytes (estimated)</td>
</tr>
<tr>
<td>approx_free_percent</td>
<td>float8</td>
<td>Percentage of free space</td>
</tr>
</tbody>
</table>

In the above output, the free space figures may not match the pgstattuple output exactly, because the free space map gives us an exact figure, but is not guaranteed to be accurate to the byte.

**F.30.2. Authors**

Tatsuo Ishii, Satoshi Nagayasu and Abhijit Menon-Sen

**F.31. pg_trgm**

The pg_trgm module provides functions and operators for determining the similarity of alphanumeric text based on trigram matching, as well as index operator classes that support fast searching for similar strings.
### F.31.1. Trigram (or Trigraph) Concepts

A trigram is a group of three consecutive characters taken from a string. We can measure the similarity of two strings by counting the number of trigrams they share. This simple idea turns out to be very effective for measuring the similarity of words in many natural languages.

**Note**

`pg_trgm` ignores non-word characters (non-alphanumerics) when extracting trigrams from a string. Each word is considered to have two spaces prefixed and one space suffixed when determining the set of trigrams contained in the string. For example, the set of trigrams in the string “cat” is “c”, “ca”, “cat”, and “at ”. The set of trigrams in the string “foo|bar” is “f”, “fo”, “foo”, “oo “, “ b”, “ ba”, “bar”, and “ar ”.

### F.31.2. Functions and Operators

The functions provided by the `pg_trgm` module are shown in Table F.24, the operators in Table F.25.

**Table F.24. pg_trgm Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>similarity(text, text)</td>
<td>real</td>
<td>Returns a number that indicates how similar the two arguments are. The range of the result is zero (indicating that the two strings are completely dissimilar) to one (indicating that the two strings are identical).</td>
</tr>
<tr>
<td>show_trgm(text)</td>
<td>text[]</td>
<td>Returns an array of all the trigrams in the given string. (In practice this is seldom useful except for debugging.)</td>
</tr>
<tr>
<td>word_similarity(text, text)</td>
<td>real</td>
<td>Returns a number that indicates the greatest similarity between the set of trigrams in the first string and any continuous extent of an ordered set of trigrams in the second string. For details, see the explanation below.</td>
</tr>
<tr>
<td>strict_word_similarity(text, text)</td>
<td>real</td>
<td>Same as <code>word_similarity(text, text)</code>, but forces extent boundaries to match word boundaries. Since we don't have cross-word trigrams, this function actually returns greatest similarity between first string and any continuous extent of words of the second string.</td>
</tr>
<tr>
<td>show_limit()</td>
<td>real</td>
<td>Returns the current similarity threshold used by the % operator. This sets the minimum similarity between two words for them to be considered similar enough</td>
</tr>
</tbody>
</table>
### Function
<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>set_limit(real)</td>
<td>real</td>
<td>Sets the current similarity threshold that is used by the % operator. The threshold must be between 0 and 1 (default is 0.3). Returns the same value passed in (deprecated).</td>
</tr>
</tbody>
</table>

Consider the following example:

```sql
# SELECT word_similarity('word', 'two words');
word_similarity
-----------------
0.8
(1 row)
```

In the first string, the set of trigrams is `{"w","wo","wor","ord","rd "}`. In the second string, the ordered set of trigrams is `{"t","tw","two","wo ","w","wo","wor","ord","rds","ds "}`. The most similar extent of an ordered set of trigrams in the second string is `{"w","wo","wor","ord"}`, and the similarity is 0.8.

This function returns a value that can be approximately understood as the greatest similarity between the first string and any substring of the second string. However, this function does not add padding to the boundaries of the extent. Thus, the number of additional characters present in the second string is not considered, except for the mismatched word boundaries.

At the same time, `strict_word_similarity(text, text)` selects an extent of words in the second string. In the example above, `strict_word_similarity(text, text)` would select the extent of a single word 'words', whose set of trigrams is `{"w","wo","wor","ord","rds","ds "}`.

```sql
# SELECT strict_word_similarity('word', 'two words'),
similarity('word', 'words');
strict_word_similarity | similarity
------------------------+------------
0.571429 | 0.571429
(1 row)
```

Thus, the `strict_word_similarity(text, text)` function is useful for finding the similarity to whole words, while `word_similarity(text, text)` is more suitable for finding the similarity for parts of words.

### Table F.25. pg_trgm Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>text % text</td>
<td>boolean</td>
<td>Returns true if its arguments have a similarity that is greater than the current similarity threshold set by pg_trgm.similarity_threshold.</td>
</tr>
<tr>
<td>text &lt; text</td>
<td>boolean</td>
<td>Returns true if the similarity between the trigram set in the first argument and a contin-</td>
</tr>
</tbody>
</table>
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>text %&gt; text</code></td>
<td>boolean</td>
<td>Commutator of the <code>&lt;%</code> operator.</td>
</tr>
<tr>
<td><code>text &lt;&lt;&lt; text</code></td>
<td>boolean</td>
<td>Returns <code>true</code> if its second argument has a continuous extent of an ordered trigram set that matches word boundaries, and its similarity to the trigram set of the first argument is greater than the current strict word similarity threshold set by the <code>pg_trgm.strict_word_similarity_threshold</code> parameter.</td>
</tr>
<tr>
<td><code>text %&gt;&gt; text</code></td>
<td>boolean</td>
<td>Commutator of the <code>&lt;&lt;%</code> operator.</td>
</tr>
<tr>
<td><code>text &lt;-&gt; text</code></td>
<td>real</td>
<td>Returns the “distance” between the arguments, that is one minus the <code>similarity()</code> value.</td>
</tr>
<tr>
<td><code>text &lt;&lt;- text</code></td>
<td>real</td>
<td>Returns the “distance” between the arguments, that is one minus the <code>word_similarity()</code> value.</td>
</tr>
<tr>
<td><code>text &lt;-&gt;&gt; text</code></td>
<td>real</td>
<td>Commutator of the <code>&lt;&lt;-&gt;</code> operator.</td>
</tr>
<tr>
<td><code>text &lt;&lt;&lt;-&gt; text</code></td>
<td>real</td>
<td>Returns the “distance” between the arguments, that is one minus the <code>strict_word_similarity()</code> value.</td>
</tr>
<tr>
<td><code>text &lt;&lt;-&gt;&gt; text</code></td>
<td>real</td>
<td>Commutator of the <code>&lt;&lt;&lt;-&gt;</code> operator.</td>
</tr>
</tbody>
</table>

### F.31.3. GUC Parameters

- **`pg_trgm.similarity_threshold(real)`**
  
  Sets the current similarity threshold that is used by the `%` operator. The threshold must be between 0 and 1 (default is 0.3).

- **`pg_trgm.word_similarity_threshold(real)`**
  
  Sets the current word similarity threshold that is used by `<%` and `%>` operators. The threshold must be between 0 and 1 (default is 0.6).

### F.31.4. Index Support

The `pg_trgm` module provides GiST and GIN index operator classes that allow you to create an index over a text column for the purpose of very fast similarity searches. These index types support the above-described similarity operators, and additionally support trigram-based index searches for `LIKE`. 
Additional Supplied Modules

ILIKE, ~ and ~* queries. (These indexes do not support equality nor simple comparison operators, so you may need a regular B-tree index too.)

Example:

```
CREATE TABLE test_trgm (t text);
CREATE INDEX trgm_idx ON test_trgm USING GIST (t gist_trgm_ops);
```

or

```
CREATE INDEX trgm_idx ON test_trgm USING GIN (t gin_trgm_ops);
```

At this point, you will have an index on the \( t \) column that you can use for similarity searching. A typical query is

```
SELECT t, similarity(t, 'word') AS sml
   FROM test_trgm
   WHERE t % 'word'
   ORDER BY sml DESC, t;
```

This will return all values in the text column that are sufficiently similar to \( \text{word} \), sorted from best match to worst. The index will be used to make this a fast operation even over very large data sets.

A variant of the above query is

```
SELECT t, t <-> 'word' AS dist
   FROM test_trgm
   ORDER BY dist LIMIT 10;
```

This can be implemented quite efficiently by GiST indexes, but not by GIN indexes. It will usually beat the first formulation when only a small number of the closest matches is wanted.

Also you can use an index on the \( t \) column for word similarity or strict word similarity. Typical queries are:

```
SELECT t, word_similarity('word', t) AS sml
   FROM test_trgm
   WHERE 'word' <\% t
   ORDER BY sml DESC, t;
```

and

```
SELECT t, strict_word_similarity('word', t) AS sml
   FROM test_trgm
   WHERE 'word' <<\% t
   ORDER BY sml DESC, t;
```

This will return all values in the text column for which there is a continuous extent in the corresponding ordered trigram set that is sufficiently similar to the trigram set of \( \text{word} \), sorted from best match to worst. The index will be used to make this a fast operation even over very large data sets.

Possible variants of the above queries are:

```
SELECT t, 'word' <-> t AS dist
   FROM test_trgm
```

3411
ORDER BY dist LIMIT 10;

and

SELECT t, 'word' <<<-> t AS dist
FROM test_trgm
ORDER BY dist LIMIT 10;

This can be implemented quite efficiently by GiST indexes, but not by GIN indexes.

Beginning in PostgreSQL 9.1, these index types also support index searches for LIKE and ILIKE, for example

SELECT * FROM test_trgm WHERE t LIKE '%foo%bar';

The index search works by extracting trigrams from the search string and then looking these up in the index. The more trigrams in the search string, the more effective the index search is. Unlike B-tree based searches, the search string need not be left-anchored.

Beginning in PostgreSQL 9.3, these index types also support index searches for regular-expression matches (~ and ~* operators), for example

SELECT * FROM test_trgm WHERE t ~ '(foo|bar)';

The index search works by extracting trigrams from the regular expression and then looking these up in the index. The more trigrams that can be extracted from the regular expression, the more effective the index search is. Unlike B-tree based searches, the search string need not be left-anchored.

For both LIKE and regular-expression searches, keep in mind that a pattern with no extractable trigrams will degenerate to a full-index scan.

The choice between GiST and GIN indexing depends on the relative performance characteristics of GiST and GIN, which are discussed elsewhere.

### F.31.5. Text Search Integration

Trigram matching is a very useful tool when used in conjunction with a full text index. In particular it can help to recognize misspelled input words that will not be matched directly by the full text search mechanism.

The first step is to generate an auxiliary table containing all the unique words in the documents:

```
CREATE TABLE words AS SELECT word FROM
  ts_stat('SELECT to_tsvector(''simple'', bodytext) FROM
  documents');
```

where `documents` is a table that has a text field `bodytext` that we wish to search. The reason for using the `simple` configuration with the `to_tsvector` function, instead of using a language-specific configuration, is that we want a list of the original (unstemmed) words.

Next, create a trigram index on the word column:

```
CREATE INDEX words_idx ON words USING GIN (word gin_trgm_ops);
```

Now, a `SELECT` query similar to the previous example can be used to suggest spellings for misspelled words in user search terms. A useful extra test is to require that the selected words are also of similar length to the misspelled word.
F.31.6. References

GiST Development Site http://www.sai.msu.su/~megera/postgres/gist/

Tsearch2 Development Site http://www.sai.msu.su/~megera/postgres/gist/tsearch/V2/

F.31.7. Authors

Oleg Bartunov <oleg@sai.msu.su>, Moscow, Moscow University, Russia

Teodor Sigaev <teodor@sigaev.ru>, Moscow, Delta-Soft Ltd., Russia

Alexander Korotkov <a.korotkov@postgrespro.ru>, Moscow, Postgres Professional, Russia

Documentation: Christopher Kings-Lynne

This module is sponsored by Delta-Soft Ltd., Moscow, Russia.

F.32. pg_visibility

The pg_visibility module provides a means for examining the visibility map (VM) and page-level visibility information of a table. It also provides functions to check the integrity of a visibility map and to force it to be rebuilt.

Three different bits are used to store information about page-level visibility. The all-visible bit in the visibility map indicates that every tuple in the corresponding page of the relation is visible to every current and future transaction. The all-frozen bit in the visibility map indicates that every tuple in the page is frozen; that is, no future vacuum will need to modify the page until such time as a tuple is inserted, updated, deleted, or locked on that page. The page header's PD_ALL_VISIBLE bit has the same meaning as the all-visible bit in the visibility map, but is stored within the data page itself rather than in a separate data structure. These two bits will normally agree, but the page's all-visible bit can sometimes be set while the visibility map bit is clear after a crash recovery. The reported values can also disagree because of a change that occurs after pg_visibility examines the visibility map and before it examines the data page. Any event that causes data corruption can also cause these bits to disagree.

Functions that display information about PD_ALL_VISIBLE bits are much more costly than those that only consult the visibility map, because they must read the relation's data blocks rather than only the (much smaller) visibility map. Functions that check the relation's data blocks are similarly expensive.

F.32.1. Functions

\texttt{pg_visibility\_map(relation regclass, blkno bigint, all\_visible OUT boolean, all\_frozen OUT boolean) returns record}

Returns the all-visible and all-frozen bits in the visibility map for the given block of the given relation.
pg_visibility(relation regclass, blkno bigint, all_visible OUT boolean, all_frozen OUT boolean, pd_all_visible OUT boolean) returns record

Returns the all-visible and all-frozen bits in the visibility map for the given block of the given relation, plus the PD_ALL_VISIBLE bit of that block.

pg_visibility_map(relation regclass, blkno OUT bigint, all_visible OUT boolean, all_frozen OUT boolean) returns setof record

Returns the all-visible and all-frozen bits in the visibility map for each block of the given relation.

pg_visibility(relation regclass, blkno OUT bigint, all_visible OUT boolean, all_frozen OUT boolean, pd_all_visible OUT boolean) returns setof record

Returns the all-visible and all-frozen bits in the visibility map for each block of the given relation, plus the PD_ALL_VISIBLE bit of each block.

pg_visibility_map_summary(relation regclass, all_visible OUT bigint, all_frozen OUT bigint) returns record

Returns the number of all-visible pages and the number of all-frozen pages in the relation according to the visibility map.

pg_check_frozen(relation regclass, t_ctid OUT tid) returns setof tid

Returns the TIDs of non-frozen tuples stored in pages marked all-frozen in the visibility map. If this function returns a non-empty set of TIDs, the visibility map is corrupt.

pg_check_visible(relation regclass, t_ctid OUT tid) returns setof tid

Returns the TIDs of non-all-visible tuples stored in pages marked all-visible in the visibility map. If this function returns a non-empty set of TIDs, the visibility map is corrupt.

pg_truncate_visibility_map(relation regclass) returns void

Truncates the visibility map for the given relation. This function is useful if you believe that the visibility map for the relation is corrupt and wish to force rebuilding it. The first VACUUM executed on the given relation after this function is executed will scan every page in the relation and rebuild the visibility map. (Until that is done, queries will treat the visibility map as containing all zeroes.)

By default, these functions are executable only by superusers and members of the pg_stat_scan_tables role, with the exception of pg_truncate_visibility_map(relation regclass) which can only be executed by superusers.

F.32.2. Author

Robert Haas <rhaas@postgresql.org>

F.33. postgres_fdw

The postgres_fdw module provides the foreign-data wrapper postgres_fdw, which can be used to access data stored in external PostgreSQL servers.

The functionality provided by this module overlaps substantially with the functionality of the older dblink module. But postgres_fdw provides more transparent and standards-compliant syntax for accessing remote tables, and can give better performance in many cases.

To prepare for remote access using postgres_fdw:
1. Install the `postgres_fdw` extension using `CREATE EXTENSION`.
2. Create a foreign server object, using `CREATE SERVER`, to represent each remote database you want to connect to. Specify connection information, except `user` and `password`, as options of the server object.
3. Create a user mapping, using `CREATE USER MAPPING`, for each database user you want to allow to access each foreign server. Specify the remote user name and password to use as `user` and `password` options of the user mapping.
4. Create a foreign table, using `CREATE FOREIGN TABLE` or `IMPORT FOREIGN SCHEMA`, for each remote table you want to access. The columns of the foreign table must match the referenced remote table. You can, however, use table and/or column names different from the remote table's, if you specify the correct remote names as options of the foreign table object.

Now you need only `SELECT` from a foreign table to access the data stored in its underlying remote table. You can also modify the remote table using `INSERT`, `UPDATE`, or `DELETE`. (Of course, the remote user you have specified in your user mapping must have privileges to do these things.)

Note that `postgres_fdw` currently lacks support for `INSERT` statements with an `ON CONFLICT DO UPDATE` clause. However, the `ON CONFLICT DO NOTHING` clause is supported, provided a unique index inference specification is omitted.

It is generally recommended that the columns of a foreign table be declared with exactly the same data types, and collations if applicable, as the referenced columns of the remote table. Although `postgres_fdw` is currently rather forgiving about performing data type conversions at need, surprising semantic anomalies may arise when types or collations do not match, due to the remote server interpreting `WHERE` clauses slightly differently from the local server.

Note that a foreign table can be declared with fewer columns, or with a different column order, than its underlying remote table has. Matching of columns to the remote table is by name, not position.

### F.33.1. FDW Options of `postgres_fdw`

#### F.33.1.1. Connection Options

A foreign server using the `postgres_fdw` foreign data wrapper can have the same options that `libpq` accepts in connection strings, as described in Section 34.1.2, except that these options are not allowed:

- `user` and `password` (specify these in a user mapping, instead)
- `client_encoding` (this is automatically set from the local server encoding)
- `fallback_application_name` (always set to `postgres_fdw`)

Only superusers may connect to foreign servers without password authentication, so always specify the `password` option for user mappings belonging to non-superusers.

#### F.33.1.2. Object Name Options

These options can be used to control the names used in SQL statements sent to the remote PostgreSQL server. These options are needed when a foreign table is created with names different from the underlying remote table's names.

- `schema_name`
  
  This option, which can be specified for a foreign table, gives the schema name to use for the foreign table on the remote server. If this option is omitted, the name of the foreign table's schema is used.

- `table_name`
  
  This option, which can be specified for a foreign table, gives the table name to use for the foreign table on the remote server. If this option is omitted, the foreign table's name is used.
Additional Supplied Modules

F.33.1.3. Cost Estimation Options

`postgres_fdw` retrieves remote data by executing queries against remote servers, so ideally the estimated cost of scanning a foreign table should be whatever it costs to be done on the remote server, plus some overhead for communication. The most reliable way to get such an estimate is to ask the remote server and then add something for overhead — but for simple queries, it may not be worth the cost of an additional remote query to get a cost estimate. So `postgres_fdw` provides the following options to control how cost estimation is done:

**use_remote_estimate**

This option, which can be specified for a foreign table or a foreign server, controls whether `postgres_fdw` issues remote EXPLAIN commands to obtain cost estimates. A setting for a foreign table overrides any setting for its server, but only for that table. The default is `false`.

**fdw_startup_cost**

This option, which can be specified for a foreign server, is a numeric value that is added to the estimated startup cost of any foreign-table scan on that server. This represents the additional overhead of establishing a connection, parsing and planning the query on the remote side, etc. The default value is `100`.

**fdw_tuple_cost**

This option, which can be specified for a foreign server, is a numeric value that is used as extra cost per-tuple for foreign-table scans on that server. This represents the additional overhead of data transfer between servers. You might increase or decrease this number to reflect higher or lower network delay to the remote server. The default value is `0.01`.

When `use_remote_estimate` is true, `postgres_fdw` obtains row count and cost estimates from the remote server and then adds `fdw_startup_cost` and `fdw_tuple_cost` to the cost estimates. When `use_remote_estimate` is false, `postgres_fdw` performs local row count and cost estimation and then adds `fdw_startup_cost` and `fdw_tuple_cost` to the cost estimates. This local estimation is unlikely to be very accurate unless local copies of the remote table's statistics are available. Running `ANALYZE` on the foreign table is the way to update the local statistics; this will perform a scan of the remote table and then calculate and store statistics just as though the table were local. Keeping local statistics can be a useful way to reduce per-query planning overhead for a remote table — but if the remote table is frequently updated, the local statistics will soon be obsolete.

F.33.1.4. Remote Execution Options

By default, only `WHERE` clauses using built-in operators and functions will be considered for execution on the remote server. Clauses involving non-built-in functions are checked locally after rows are fetched. If such functions are available on the remote server and can be relied on to produce the same results as they do locally, performance can be improved by sending such `WHERE` clauses for remote execution. This behavior can be controlled using the following option:

**extensions**

This option is a comma-separated list of names of PostgreSQL extensions that are installed, in compatible versions, on both the local and remote servers. Functions and operators that are immutable and belong to a listed extension will be considered shippable to the remote server. This option can only be specified for foreign servers, not per-table.
When using the extensions option, it is the user's responsibility that the listed extensions exist and behave identically on both the local and remote servers. Otherwise, remote queries may fail or behave unexpectedly.

**fetch_size**

This option specifies the number of rows `postgres_fdw` should get in each fetch operation. It can be specified for a foreign table or a foreign server. The option specified on a table overrides an option specified for the server. The default is 100.

### F.33.1.5. Updatability Options

By default all foreign tables using `postgres_fdw` are assumed to be updatable. This may be overridden using the following option:

**updatable**

This option controls whether `postgres_fdw` allows foreign tables to be modified using `INSERT`, `UPDATE` and `DELETE` commands. It can be specified for a foreign table or a foreign server. A table-level option overrides a server-level option. The default is `true`.

Of course, if the remote table is not in fact updatable, an error would occur anyway. Use of this option primarily allows the error to be thrown locally without querying the remote server. Note however that the `information_schema` views will report a `postgres_fdw` foreign table to be updatable (or not) according to the setting of this option, without any check of the remote server.

### F.33.1.6. Importing Options

`postgres_fdw` is able to import foreign table definitions using `IMPORT FOREIGN SCHEMA`. This command creates foreign table definitions on the local server that match tables or views present on the remote server. If the remote tables to be imported have columns of user-defined data types, the local server must have compatible types of the same names.

Importing behavior can be customized with the following options (given in the `IMPORT FOREIGN SCHEMA` command):

**import_collate**

This option controls whether column `COLLATE` options are included in the definitions of foreign tables imported from a foreign server. The default is `true`. You might need to turn this off if the remote server has a different set of collation names than the local server does, which is likely to be the case if it's running on a different operating system.

**import_default**

This option controls whether column `DEFAULT` expressions are included in the definitions of foreign tables imported from a foreign server. The default is `false`. If you enable this option, be wary of defaults that might get computed differently on the local server than they would be on the remote server; `nextval()` is a common source of problems. The `IMPORT` will fail altogether if an imported default expression uses a function or operator that does not exist locally.

**import_not_null**

This option controls whether column `NOT NULL` constraints are included in the definitions of foreign tables imported from a foreign server. The default is `true`.

Note that constraints other than `NOT NULL` will never be imported from the remote tables. Although PostgreSQL does support `CHECK` constraints on foreign tables, there is no provision for importing
them automatically, because of the risk that a constraint expression could evaluate differently on the local and remote servers. Any such inconsistency in the behavior of a CHECK constraint could lead to hard-to-detect errors in query optimization. So if you wish to import CHECK constraints, you must do so manually, and you should verify the semantics of each one carefully. For more detail about the treatment of CHECK constraints on foreign tables, see CREATE FOREIGN TABLE.

Tables or foreign tables which are partitions of some other table are automatically excluded. Partitioned tables are imported, unless they are a partition of some other table. Since all data can be accessed through the partitioned table which is the root of the partitioning hierarchy, this approach should allow access to all the data without creating extra objects.

F.33.2. Connection Management

postgres_fdw establishes a connection to a foreign server during the first query that uses a foreign table associated with the foreign server. This connection is kept and re-used for subsequent queries in the same session. However, if multiple user identities (user mappings) are used to access the foreign server, a connection is established for each user mapping.

F.33.3. Transaction Management

During a query that references any remote tables on a foreign server, postgres_fdw opens a transaction on the remote server if one is not already open corresponding to the current local transaction. The remote transaction is committed or aborted when the local transaction commits or aborts. Savepoints are similarly managed by creating corresponding remote savepoints.

The remote transaction uses SERIALizable isolation level when the local transaction has SERIALizable isolation level; otherwise it uses REPEATable READ isolation level. This choice ensures that if a query performs multiple table scans on the remote server, it will get snapshot-consistent results for all the scans. A consequence is that successive queries within a single transaction will see the same data from the remote server, even if concurrent updates are occurring on the remote server due to other activities. That behavior would be expected anyway if the local transaction uses SERIALizable or REPEATable READ isolation level, but it might be surprising for a READ COMMITTED local transaction. A future PostgreSQL release might modify these rules.

F.33.4. Remote Query Optimization

postgres_fdw attempts to optimize remote queries to reduce the amount of data transferred from foreign servers. This is done by sending query WHERE clauses to the remote server for execution, and by not retrieving table columns that are not needed for the current query. To reduce the risk of misexecution of queries, WHERE clauses are not sent to the remote server unless they use only data types, operators, and functions that are built-in or belong to an extension that's listed in the foreign server's extensions option. Operators and functions in such clauses must be IMMUTABLE as well. For an UPDATE or DELETE query, postgres_fdw attempts to optimize the query execution by sending the whole query to the remote server if there are no query WHERE clauses that cannot be sent to the remote server, no local joins for the query, no row-level local BEFORE or AFTER triggers on the target table, and no CHECK OPTION constraints from parent views. In UPDATE, expressions to assign to target columns must use only built-in data types, IMMUTABLE operators, or IMMUTABLE functions, to reduce the risk of misexecution of the query.

When postgres_fdw encounters a join between foreign tables on the same foreign server, it sends the entire join to the foreign server, unless for some reason it believes that it will be more efficient to fetch rows from each table individually, or unless the table references involved are subject to different user mappings. While sending the JOIN clauses, it takes the same precautions as mentioned above for the WHERE clauses.

The query that is actually sent to the remote server for execution can be examined using EXPLAIN VERBOSE.
F.33.5. Remote Query Execution Environment

In the remote sessions opened by `postgres_fdw`, the `search_path` parameter is set to just `pg_catalog`, so that only built-in objects are visible without schema qualification. This is not an issue for queries generated by `postgres_fdw` itself, because it always supplies such qualification. However, this can pose a hazard for functions that are executed on the remote server via triggers or rules on remote tables. For example, if a remote table is actually a view, any functions used in that view will be executed with the restricted search path. It is recommended to schema-qualify all names in such functions, or else attach `SET search_path` options (see `CREATE FUNCTION`) to such functions to establish their expected search path environment.

`postgres_fdw` likewise establishes remote session settings for various parameters:

- `TimeZone` is set to `UTC`
- `DateStyle` is set to `ISO`
- `IntervalStyle` is set to `postgres`
- `extra_float_digits` is set to `3` for remote servers 9.0 and newer and is set to `2` for older versions

These are less likely to be problematic than `search_path`, but can be handled with function `SET` options if the need arises.

It is not recommended that you override this behavior by changing the session-level settings of these parameters; that is likely to cause `postgres_fdw` to malfunction.

F.33.6. Cross-Version Compatibility

`postgres_fdw` can be used with remote servers dating back to PostgreSQL 8.3. Read-only capability is available back to 8.1. A limitation however is that `postgres_fdw` generally assumes that immutable built-in functions and operators are safe to send to the remote server for execution, if they appear in a `WHERE` clause for a foreign table. Thus, a built-in function that was added since the remote server's release might be sent to it for execution, resulting in “function does not exist” or a similar error. This type of failure can be worked around by rewriting the query, for example by embedding the foreign table reference in a sub-select with `OFFSET 0` as an optimization fence, and placing the problematic function or operator outside the sub-select.

F.33.7. Examples

Here is an example of creating a foreign table with `postgres_fdw`. First install the extension:

```
CREATE EXTENSION postgres_fdw;
```

Then create a foreign server using `CREATE SERVER`. In this example we wish to connect to a PostgreSQL server on host `192.83.123.89` listening on port `5432`. The database to which the connection is made is named `foreign_db` on the remote server:

```
CREATE SERVER foreign_server
    FOREIGN DATA WRAPPER postgres_fdw
    OPTIONS (host '192.83.123.89', port '5432', dbname 'foreign_db');
```

A user mapping, defined with `CREATE USER MAPPING`, is needed as well to identify the role that will be used on the remote server:

```
CREATE USER MAPPING FOR local_user
    SERVER foreign_server
```
OPTIONS (user 'foreign_user', password 'password');

Now it is possible to create a foreign table with CREATE FOREIGN TABLE. In this example we wish to access the table named some_schema.some_table on the remote server. The local name for it will be foreign_table:

```sql
CREATE FOREIGN TABLE foreign_table (
    id integer NOT NULL,
    data text
) SERVER foreign_server
OPTIONS (schema_name 'some_schema', table_name 'some_table');
```

It's essential that the data types and other properties of the columns declared in CREATE FOREIGN TABLE match the actual remote table. Column names must match as well, unless you attach column_name options to the individual columns to show how they are named in the remote table. In many cases, use of IMPORT FOREIGN SCHEMA is preferable to constructing foreign table definitions manually.

F.33.8. Author

Shigeru Hanada <shigeru.hanada@gmail.com>

F.34. seg

This module implements a data type seg for representing line segments, or floating point intervals. seg can represent uncertainty in the interval endpoints, making it especially useful for representing laboratory measurements.

F.34.1. Rationale

The geometry of measurements is usually more complex than that of a point in a numeric continuum. A measurement is usually a segment of that continuum with somewhat fuzzy limits. The measurements come out as intervals because of uncertainty and randomness, as well as because the value being measured may naturally be an interval indicating some condition, such as the temperature range of stability of a protein.

Using just common sense, it appears more convenient to store such data as intervals, rather than pairs of numbers. In practice, it even turns out more efficient in most applications.

Further along the line of common sense, the fuzziness of the limits suggests that the use of traditional numeric data types leads to a certain loss of information. Consider this: your instrument reads 6.50, and you input this reading into the database. What do you get when you fetch it? Watch:

```sql
test=> select 6.50 :: float8 as "pH";
     pH
---
    6.5
(1 row)
```

In the world of measurements, 6.50 is not the same as 6.5. It may sometimes be critically different. The experimenters usually write down (and publish) the digits they trust. 6.50 is actually a fuzzy interval contained within a bigger and even fuzzier interval, 6.5, with their center points being (probably) the only common feature they share. We definitely do not want such different data items to appear the same.
Conclusion? It is nice to have a special data type that can record the limits of an interval with arbitrarily variable precision. Variable in the sense that each data element records its own precision.

Check this out:

```sql
test=> select '6.25 .. 6.50 '::seg as "pH";
pH
------------
 6.25 .. 6.50
(1 row)
```

**F.34.2. Syntax**

The external representation of an interval is formed using one or two floating-point numbers joined by the range operator (.. or ...). Alternatively, it can be specified as a center point plus or minus a deviation. Optional certainty indicators (<, > or ~) can be stored as well. (Certainty indicators are ignored by all the built-in operators, however.) Table F.26 gives an overview of allowed representations; Table F.27 shows some examples.

In Table F.26, x, y, and delta denote floating-point numbers. x and y, but not delta, can be preceded by a certainty indicator.

**Table F.26. seg External Representations**

<table>
<thead>
<tr>
<th>Representation</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>x</td>
<td>Single value (zero-length interval)</td>
</tr>
<tr>
<td>x .. y</td>
<td>Interval from x to y</td>
</tr>
<tr>
<td>x (+-) delta</td>
<td>Interval from x - delta to x + delta</td>
</tr>
<tr>
<td>x ..</td>
<td>Open interval with lower bound x</td>
</tr>
<tr>
<td>.. x</td>
<td>Open interval with upper bound x</td>
</tr>
</tbody>
</table>

**Table F.27. Examples of Valid seg Input**

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>5.0</td>
<td>Creates a zero-length segment (a point, if you will)</td>
</tr>
<tr>
<td>~5.0</td>
<td>Creates a zero-length segment and records ~ in the data. ~ is ignored by seg operations, but is preserved as a comment.</td>
</tr>
<tr>
<td>&lt;5.0</td>
<td>Creates a point at 5.0. &lt; is ignored but is preserved as a comment.</td>
</tr>
<tr>
<td>&gt;5.0</td>
<td>Creates a point at 5.0. &gt; is ignored but is preserved as a comment.</td>
</tr>
<tr>
<td>5(+-)0.3</td>
<td>Creates an interval 4.7 .. 5.3. Note that the (+-) notation isn't preserved.</td>
</tr>
<tr>
<td>50 ..</td>
<td>Everything that is greater than or equal to 50</td>
</tr>
<tr>
<td>.. 0</td>
<td>Everything that is less than or equal to 0</td>
</tr>
<tr>
<td>1.5e-2 .. 2E-2</td>
<td>Creates an interval 0.015 .. 0.02</td>
</tr>
<tr>
<td>1 ... 2</td>
<td>The same as 1 ... 2, or 1 .. 2, or 1 .. 2 (spaces around the range operator are ignored)</td>
</tr>
</tbody>
</table>

Because ... is widely used in data sources, it is allowed as an alternative spelling of ... Unfortunately, this creates a parsing ambiguity: it is not clear whether the upper bound in 0 ... 23 is meant to be 23 or 0.23. This is resolved by requiring at least one digit before the decimal point in all numbers in seg input.
As a sanity check, \texttt{seg} rejects intervals with the lower bound greater than the upper, for example 5 .. 2.

### F.34.3. Precision

\texttt{seg} values are stored internally as pairs of 32-bit floating point numbers. This means that numbers with more than 7 significant digits will be truncated.

Numbers with 7 or fewer significant digits retain their original precision. That is, if your query returns 0.00, you will be sure that the trailing zeroes are not the artifacts of formatting: they reflect the precision of the original data. The number of leading zeroes does not affect precision: the value 0.0067 is considered to have just 2 significant digits.

### F.34.4. Usage

The \texttt{seg} module includes a GiST index operator class for \texttt{seg} values. The operators supported by the GiST operator class are shown in Table F.28.

#### Table F.28. Seg GiST Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>([a, b] &lt;&lt; [c, d])</td>
<td>([a, b] ) is entirely to the left of ([c, d]). That is, ([a, b] &lt;&lt; [c, d]) is true if (b &lt; c) and false otherwise.</td>
</tr>
<tr>
<td>([a, b] &gt;&gt; [c, d])</td>
<td>([a, b] ) is entirely to the right of ([c, d]). That is, ([a, b] &gt;&gt; [c, d]) is true if (a &gt; d) and false otherwise.</td>
</tr>
<tr>
<td>([a, b] &amp;&lt; [c, d])</td>
<td>Overlaps or is left of — This might be better read as “does not extend to right of”. It is true when (b \leq d).</td>
</tr>
<tr>
<td>([a, b] &amp;&gt; [c, d])</td>
<td>Overlaps or is right of — This might be better read as “does not extend to left of”. It is true when (a \geq c).</td>
</tr>
<tr>
<td>([a, b] = [c, d])</td>
<td>Same as — The segments ([a, b]) and ([c, d]) are identical, that is, (a = c) and (b = d).</td>
</tr>
<tr>
<td>([a, b] &amp;&amp; [c, d])</td>
<td>The segments ([a, b]) and ([c, d]) overlap.</td>
</tr>
<tr>
<td>([a, b] @&gt; [c, d])</td>
<td>The segment ([a, b]) contains the segment ([c, d]), that is, (a \leq c) and (b \geq d).</td>
</tr>
<tr>
<td>([a, b] &lt;@ [c, d])</td>
<td>The segment ([a, b]) is contained in ([c, d]), that is, (a \geq c) and (b \leq d).</td>
</tr>
</tbody>
</table>

(Before PostgreSQL 8.2, the containment operators \(@>\) and \(<@\) were respectively called \(\&\&\) and \(~\). These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The standard B-tree operators are also provided, for example

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>([a, b] &lt; [c, d])</td>
<td>Less than</td>
</tr>
<tr>
<td>([a, b] &gt; [c, d])</td>
<td>Greater than</td>
</tr>
</tbody>
</table>

These operators do not make a lot of sense for any practical purpose but sorting. These operators first compare \((a)\) to \((c)\), and if these are equal, compare \((b)\) to \((d)\). That results in reasonably good sorting in most cases, which is useful if you want to use ORDER BY with this type.
F.34.5. Notes

For examples of usage, see the regression test `sql/seg.sql`.

The mechanism that converts `+-` to regular ranges isn't completely accurate in determining the number of significant digits for the boundaries. For example, it adds an extra digit to the lower boundary if the resulting interval includes a power of ten:

```sql
postgres=> select '10(+-)1':::seg as seg;
  seg
---------
 9.0 .. 11             -- should be: 9 .. 11
```

The performance of an R-tree index can largely depend on the initial order of input values. It may be very helpful to sort the input table on the `seg` column; see the script `sort-segments.pl` for an example.

F.34.6. Credits

Original author: Gene Selkov, Jr. <selkovjr@mcs.anl.gov>, Mathematics and Computer Science Division, Argonne National Laboratory.

My thanks are primarily to Prof. Joe Hellerstein (http://db.cs.berkeley.edu/jmh/) for elucidating the gist of the GiST (http://gist.cs.berkeley.edu/). I am also grateful to all Postgres developers, present and past, for enabling myself to create my own world and live undisturbed in it. And I would like to acknowledge my gratitude to Argonne Lab and to the U.S. Department of Energy for the years of faithful support of my database research.

F.35. sepgsql

`sepgsql` is a loadable module that supports label-based mandatory access control (MAC) based on SELinux security policy.

```
Warning
The current implementation has significant limitations, and does not enforce mandatory access control for all actions. See Section F.35.7.
```

F.35.1. Overview

This module integrates with SELinux to provide an additional layer of security checking above and beyond what is normally provided by PostgreSQL. From the perspective of SELinux, this module allows PostgreSQL to function as a user-space object manager. Each table or function access initiated by a DML query will be checked against the system security policy. This check is in addition to the usual SQL permissions checking performed by PostgreSQL.

SELinux access control decisions are made using security labels, which are represented by strings such as `system_u:object_r:sepgsql_table_t:s0`. Each access control decision involves two labels: the label of the subject attempting to perform the action, and the label of the object on which the operation is to be performed. Since these labels can be applied to any sort of object, access control decisions for objects stored within the database can be (and, with this module, are) subjected to the same general criteria used for objects of any other type, such as files. This design is intended to allow a centralized security policy to protect information assets independent of the particulars of how those assets are stored.
The SECURITY LABEL statement allows assignment of a security label to a database object.

**F.35.2. Installation**

`sepgsql` can only be used on Linux 2.6.28 or higher with SELinux enabled. It is not available on any other platform. You will also need libselinux 2.1.10 or higher and selinux-policy 3.9.13 or higher (although some distributions may backport the necessary rules into older policy versions).

The `sestatus` command allows you to check the status of SELinux. A typical display is:

```
$ sestatus
SELinux status: enabled
SELinuxfs mount: /selinux
Current mode: enforcing
Mode from config file: enforcing
Policy version: 24
Policy from config file: targeted
```

If SELinux is disabled or not installed, you must set that product up first before installing this module.

To build this module, include the option `--with-selinux` in your PostgreSQL `configure` command. Be sure that the `libselinux-devel` RPM is installed at build time.

To use this module, you must include `sepgsql` in the `shared_preload_libraries` parameter in `postgresql.conf`. The module will not function correctly if loaded in any other manner. Once the module is loaded, you should execute `sepgsql.sql` in each database. This will install functions needed for security label management, and assign initial security labels.

Here is an example showing how to initialize a fresh database cluster with `sepgsql` functions and security labels installed. Adjust the paths shown as appropriate for your installation:

```
$ export PGDATA=/path/to/data/directory
$ initdb
$ vi $PGDATA/postgresql.conf
change
    #shared_preload_libraries = ''  # change
requires restart)
to
    shared_preload_libraries = 'sepgsql'  # change
requires restart)
$ for DBNAME in template0 template1 postgres; do
    postgres --single -F -c exit_on_error=true $DBNAME \
        </usr/local/pgsql/share/contrib/sepgsql.sql >/dev/null
done
```

Please note that you may see some or all of the following notifications depending on the particular versions you have of libselinux and selinux-policy:

```
/etc/selinux/targeted/contexts/sepgsql_contexts:  line 33 has invalid object type db_blobs
/etc/selinux/targeted/contexts/sepgsql_contexts:  line 36 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts:  line 37 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts:  line 38 has invalid object type db_language
```

3424
Additional Supplied Modules

/etc/selinux/targeted/contexts/sepgsql_contexts: line 39 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts: line 40 has invalid object type db_language

These messages are harmless and should be ignored.

If the installation process completes without error, you can now start the server normally.

F.35.3. Regression Tests

Due to the nature of SELinux, running the regression tests for sepgsql requires several extra configuration steps, some of which must be done as root. The regression tests will not be run by an ordinary make check or make installcheck command; you must set up the configuration and then invoke the test script manually. The tests must be run in the contrib/sepgsql directory of a configured PostgreSQL build tree. Although they require a build tree, the tests are designed to be executed against an installed server, that is they are comparable to make installcheck not make check.

First, set up sepgsql in a working database according to the instructions in Section F.35.2. Note that the current operating system user must be able to connect to the database as superuser without password authentication.

Second, build and install the policy package for the regression test. The sepgsql-regtest policy is a special purpose policy package which provides a set of rules to be allowed during the regression tests. It should be built from the policy source file sepgsql-regtest.te, which is done using make with a Makefile supplied by SELinux. You will need to locate the appropriate Makefile on your system; the path shown below is only an example. Once built, install this policy package using the semodule command, which loads supplied policy packages into the kernel. If the package is correctly installed, semodule -l should list sepgsql-regtest as an available policy package:

$ cd .../contrib/sepgsql
$ make -f /usr/share/selinux/devel/Makefile
$ sudo semodule -u sepgsql-regtest.pp
$ sudo semodule -l | grep sepgsql
sepgsql-regtest 1.07

Third, turn on sepgsql_regression_test_mode. For security reasons, the rules in sepgsql-regtest are not enabled by default; the sepgsql_regression_test_mode parameter enables the rules needed to launch the regression tests. It can be turned on using the setsebool command:

$ sudo setsebool sepgsql_regression_test_mode on
$ getsebool sepgsql_regression_test_mode
sepgsql_regression_test_mode --> on

Fourth, verify your shell is operating in the unconfined_t domain:

$ id -Z
unconfined_u:unconfined_r:unconfined_t:s0-s0:c0.c1023

See Section F.35.8 for details on adjusting your working domain, if necessary.

Finally, run the regression test script:

$ ./test_sepgsql
This script will attempt to verify that you have done all the configuration steps correctly, and then it will run the regression tests for the sepgsql module.

After completing the tests, it’s recommended you disable the sepgsql_regression_test_mode parameter:

$ sudo setsebool sepgsql_regression_test_mode off

You might prefer to remove the sepgsql-regtest policy entirely:

$ sudo semodule -r sepgsql-regtest

### F.35.4. GUC Parameters

**sepgsql.permissive**(boolean)

This parameter enables sepgsql to function in permissive mode, regardless of the system setting. The default is off. This parameter can only be set in the postgresql.conf file or on the server command line.

When this parameter is on, sepgsql functions in permissive mode, even if SELinux in general is working in enforcing mode. This parameter is primarily useful for testing purposes.

**sepgsql.debug_audit**(boolean)

This parameter enables the printing of audit messages regardless of the system policy settings. The default is off, which means that messages will be printed according to the system settings.

The security policy of SELinux also has rules to control whether or not particular accesses are logged. By default, access violations are logged, but allowed accesses are not.

This parameter forces all possible logging to be turned on, regardless of the system policy.

### F.35.5. Features

#### F.35.5.1. Controlled Object Classes

The security model of SELinux describes all the access control rules as relationships between a subject entity (typically, a client of the database) and an object entity (such as a database object), each of which is identified by a security label. If access to an unlabeled object is attempted, the object is treated as if it were assigned the label unlabeled_t.

Currently, sepgsql allows security labels to be assigned to schemas, tables, columns, sequences, views, and functions. When sepgsql is in use, security labels are automatically assigned to supported database objects at creation time. This label is called a default security label, and is decided according to the system security policy, which takes as input the creator's label, the label assigned to the new object's parent object and optionally name of the constructed object.

A new database object basically inherits the security label of the parent object, except when the security policy has special rules known as type-transition rules, in which case a different label may be applied. For schemas, the parent object is the current database; for tables, sequences, views, and functions, it is the containing schema; for columns, it is the containing table.

#### F.35.5.2. DML Permissions

For tables, db_table:select, db_table:insert, db_table:update or db_table:delete are checked for all the referenced target tables depending on the kind of statement; in addition, db_table:select is also checked for all the tables that contain columns referenced in the WHERE or RETURNING clause, as a data source for UPDATE, and so on.
Column-level permissions will also be checked for each referenced column. `db_column:select` is checked on not only the columns being read using `SELECT`, but those being referenced in other DML statements; `db_column:update` or `db_column:insert` will also be checked for columns being modified by `UPDATE` or `INSERT`.

For example, consider:

```
UPDATE t1 SET x = 2, y = func1(y) WHERE z = 100;
```

Here, `db_column:update` will be checked for `t1.x`, since it is being updated, `db_column:{select update}` will be checked for `t1.y`, since it is both updated and referenced, and `db_column:select` will be checked for `t1.z`, since it is only referenced. `db_table:{select update}` will also be checked at the table level.

For sequences, `db_sequence:get_value` is checked when we reference a sequence object using `SELECT`; however, note that we do not currently check permissions on execution of corresponding functions such as `lastval()`.

For views, `db_view:expand` will be checked, then any other required permissions will be checked on the objects being expanded from the view, individually.

For functions, `db_procedure:{execute}` will be checked when user tries to execute a function as a part of query, or using fast-path invocation. If this function is a trusted procedure, it also checks `db_procedure:{entrypoint}` permission to check whether it can perform as entry point of trusted procedure.

In order to access any schema object, `db_schema:search` permission is required on the containing schema. When an object is referenced without schema qualification, schemas on which this permission is not present will not be searched (just as if the user did not have `USAGE` privilege on the schema). If an explicit schema qualification is present, an error will occur if the user does not have the requisite permission on the named schema.

The client must be allowed to access all referenced tables and columns, even if they originated from views which were then expanded, so that we apply consistent access control rules independent of the manner in which the table contents are referenced.

The default database privilege system allows database superusers to modify system catalogs using DML commands, and reference or modify toast tables. These operations are prohibited when `sepgsql` is enabled.

### F.35.5.3. DDL Permissions

SELinux defines several permissions to control common operations for each object type; such as creation, alter, drop and relabel of security label. In addition, several object types have special permissions to control their characteristic operations; such as addition or deletion of name entries within a particular schema.

Creating a new database object requires `create` permission. SELinux will grant or deny this permission based on the client's security label and the proposed security label for the new object. In some cases, additional privileges are required:

- **CREATE DATABASE** additionally requires `getattr` permission for the source or template database.

- Creating a schema object additionally requires `add_name` permission on the parent schema.

- Creating a table additionally requires permission to create each individual table column, just as if each table column were a separate top-level object.

- Creating a function marked as LEAKPROOF additionally requires `install` permission. (This permission is also checked when LEAKPROOF is set for an existing function.)
When DROP command is executed, drop will be checked on the object being removed. Permissions will be also checked for objects dropped indirectly via CASCADE. Deletion of objects contained within a particular schema (tables, views, sequences and procedures) additionally requires remove_name on the schema.

When ALTER command is executed, setattr will be checked on the object being modified for each object types, except for subsidiary objects such as the indexes or triggers of a table, where permissions are instead checked on the parent object. In some cases, additional permissions are required:

- Moving an object to a new schema additionally requires remove_name permission on the old schema and add_name permission on the new one.
- Setting the LEAKPROOF attribute on a function requires install permission.
- Using SECURITY LABEL on an object additionally requires relabelfrom permission for the object in conjunction with its old security label and relabelto permission for the object in conjunction with its new security label. (In cases where multiple label providers are installed and the user tries to set a security label, but it is not managed by SELinux, only setattr should be checked here. This is currently not done due to implementation restrictions.)

**F.35.5.4. Trusted Procedures**

Trusted procedures are similar to security definer functions or setuid commands. SELinux provides a feature to allow trusted code to run using a security label different from that of the client, generally for the purpose of providing highly controlled access to sensitive data (e.g. rows might be omitted, or the precision of stored values might be reduced). Whether or not a function acts as a trusted procedure is controlled by its security label and the operating system security policy. For example:

```
pagres=# CREATE TABLE customer (  
    cid     int primary key,  
    cname   text,  
    credit  text  
  );
CREATE TABLE

Postgres=# SECURITY LABEL ON COLUMN customer.credit
IS 'system_u:object_r:sepgsql_secret_table_t:s0';
SECURITY LABEL
Postgres=# CREATE FUNCTION show_credit(int) RETURNS text
AS 'SELECT regexp_replace(credit, ''-[0-9]+$'', ''-xxxx'', ''g'')
FROM customer WHERE cid = $1'
LANGUAGE sql;
CREATE FUNCTION
Postgres=# SECURITY LABEL ON FUNCTION show_credit(int)
IS
'system_u:object_r:sepgsql_trusted_proc_exec_t:s0';
SECURITY LABEL
```

The above operations should be performed by an administrative user.

```
pagres=# SELECT * FROM customer;
ERROR:  SELinux: security policy violation
pagres=# SELECT cid, cname, show_credit(cid) FROM customer;
  cid |  cname  |         show_credit
---------+----------+---------------------
     1 |  taro       | 1111-2222-3333-xxxx
     2 |  hanako     | 5555-6666-7777-xxxx
(2 rows)
```
In this case, a regular user cannot reference `customer.credit` directly, but a trusted procedure `show_credit` allows the user to print the credit card numbers of customers with some of the digits masked out.

### F.35.5.5. Dynamic Domain Transitions

It is possible to use SELinux's dynamic domain transition feature to switch the security label of the client process, the client domain, to a new context, if that is allowed by the security policy. The client domain needs the `setcurrent` permission and also `dyntransition` from the old to the new domain.

Dynamic domain transitions should be considered carefully, because they allow users to switch their label, and therefore their privileges, at their option, rather than (as in the case of a trusted procedure) as mandated by the system. Thus, the `dyntransition` permission is only considered safe when used to switch to a domain with a smaller set of privileges than the original one. For example:

```sql
regression=# select sepgsql_getcon();
sepgsql_getcon
-------------------------------------------------------
unconfined_u:unconfined_r:unconfined_t:s0-s0:c0.c1023
(1 row)

regression=# SELECT
    sepgsql_setcon('unconfined_u:unconfined_r:unconfined_t:s0-s0:c1.c4');
sepgsql_setcon
----------------
t
(1 row)

regression=# SELECT
    sepgsql_setcon('unconfined_u:unconfined_r:unconfined_t:s0-s0:c1.c1023');
ERROR:  SELinux: security policy violation
```

In this example above we were allowed to switch from the larger MCS range `c1.c1023` to the smaller range `c1.c4`, but switching back was denied.

A combination of dynamic domain transition and trusted procedure enables an interesting use case that fits the typical process life-cycle of connection pooling software. Even if your connection pooling software is not allowed to run most of SQL commands, you can allow it to switch the security label of the client using the `sepgsql_setcon()` function from within a trusted procedure; that should take some credential to authorize the request to switch the client label. After that, this session will have the privileges of the target user, rather than the connection pooler. The connection pooler can later revert the security label change by again using `sepgsql_setcon()` with `NULL` argument, again invoked from within a trusted procedure with appropriate permissions checks. The point here is that only the trusted procedure actually has permission to change the effective security label, and only does so when given proper credentials. Of course, for secure operation, the credential store (table, procedure definition, or whatever) must be protected from unauthorized access.

### F.35.5.6. Miscellaneous

We reject the LOAD command across the board, because any module loaded could easily circumvent security policy enforcement.

### F.35.6. Sepgsql Functions

Table F.29 shows the available functions.
### Table F.29. Sepgsql Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sepgsql_getcon() returns text</td>
<td>Returns the client domain, the current security label of the client.</td>
</tr>
<tr>
<td>sepgsql_setcon(text) returns bool</td>
<td>Switches the client domain of the current session to the new domain, if allowed by the security policy. It also accepts NULL input as a request to transition to the client's original domain.</td>
</tr>
<tr>
<td>sepgsql_mcstrans_in(text) returns text</td>
<td>Translates the given qualified MLS/MCS range into raw format if the mcstrans daemon is running.</td>
</tr>
<tr>
<td>sepgsql_mcstrans_out(text) returns text</td>
<td>Translates the given raw MLS/MCS range into qualified format if the mcstrans daemon is running.</td>
</tr>
<tr>
<td>sepgsql_restorecon(text) returns bool</td>
<td>Sets up initial security labels for all objects within the current database. The argument may be NULL, or the name of a specfile to be used as alternative of the system default.</td>
</tr>
</tbody>
</table>

### F.35.7. Limitations

**Data Definition Language (DDL) Permissions**

Due to implementation restrictions, some DDL operations do not check permissions.

**Data Control Language (DCL) Permissions**

Due to implementation restrictions, DCL operations do not check permissions.

**Row-level access control**

PostgreSQL supports row-level access, but sepgsql does not.

**Covert channels**

sepgsql does not try to hide the existence of a certain object, even if the user is not allowed to reference it. For example, we can infer the existence of an invisible object as a result of primary key conflicts, foreign key violations, and so on, even if we cannot obtain the contents of the object. The existence of a top secret table cannot be hidden; we only hope to conceal its contents.

### F.35.8. External Resources

**SE-PostgreSQL Introduction**

This wiki page provides a brief overview, security design, architecture, administration and upcoming features.

**SELinux User's and Administrator's Guide**

This document provides a wide spectrum of knowledge to administer SELinux on your systems. It focuses primarily on Red Hat operating systems, but is not limited to them.

**Fedora SELinux FAQ**

This document answers frequently asked questions about SELinux. It focuses primarily on Fedora, but is not limited to Fedora.

---

5 https://wiki.postgresql.org/wiki/SEPostgreSQL
7 https://fedoraproject.org/wiki/SELinux_FAQ
F.36. spi

The spi module provides several workable examples of using the Server Programming Interface (SPI) and triggers. While these functions are of some value in their own right, they are even more useful as examples to modify for your own purposes. The functions are general enough to be used with any table, but you have to specify table and field names (as described below) while creating a trigger.

Each of the groups of functions described below is provided as a separately-installable extension.

F.36.1. refint — Functions for Implementing Referential Integrity

check_primary_key() and check_foreign_key() are used to check foreign key constraints. (This functionality is long since superseded by the built-in foreign key mechanism, of course, but the module is still useful as an example.)

check_primary_key() checks the referencing table. To use, create a BEFORE INSERT OR UPDATE trigger using this function on a table referencing another table. Specify as the trigger arguments: the referencing table's column name(s) which form the foreign key, the referenced table name, and the column names in the referenced table which form the primary/unique key. To handle multiple foreign keys, create a trigger for each reference.

check_foreign_key() checks the referenced table. To use, create a BEFORE DELETE OR UPDATE trigger using this function on a table referenced by other table(s). Specify as the trigger arguments: the number of referencing tables for which the function has to perform checking, the action if a referencing key is found (cascade — to delete the referencing row, restrict — to abort transaction if referencing keys exist, setnull — to set referencing key fields to null), the triggered table's column names which form the primary/unique key, then the referencing table name and column names (repeated for as many referencing tables as were specified by first argument). Note that the primary/unique key columns should be marked NOT NULL and should have a unique index.

There are examples in refint.example.

F.36.2. timetravel — Functions for Implementing Time Travel

Long ago, PostgreSQL had a built-in time travel feature that kept the insert and delete times for each tuple. This can be emulated using these functions. To use these functions, you must add to a table two columns of abstime type to store the date when a tuple was inserted (start_date) and changed/deleted (stop_date):

```
CREATE TABLE mytab (  ...
    start_date      abstime,
    stop_date       abstime
    ...
);
```

The columns can be named whatever you like, but in this discussion we'll call them start_date and stop_date.
When a new row is inserted, start_date should normally be set to current time, and stop_date to infinity. The trigger will automatically substitute these values if the inserted data contains nulls in these columns. Generally, inserting explicit non-null data in these columns should only be done when re-loading dumped data.

Tuples with stop_date equal to infinity are “valid now”, and can be modified. Tuples with a finite stop_date cannot be modified anymore — the trigger will prevent it. (If you need to do that, you can turn off time travel as shown below.)

For a modifiable row, on update only the stop_date in the tuple being updated will be changed (to current time) and a new tuple with the modified data will be inserted. Start_date in this new tuple will be set to current time and stop_date to infinity.

A delete does not actually remove the tuple but only sets its stop_date to current time.

To query for tuples “valid now”, include stop_date = 'infinity' in the query’s WHERE condition. (You might wish to incorporate that in a view.) Similarly, you can query for tuples valid at any past time with suitable conditions on start_date and stop_date.

timetavel() is the general trigger function that supports this behavior. Create a BEFORE INSERT OR UPDATE OR DELETE trigger using this function on each time-traveled table. Specify two trigger arguments: the actual names of the start_date and stop_date columns. Optionally, you can specify one to three more arguments, which must refer to columns of type text. The trigger will store the name of the current user into the first of these columns during INSERT, the second column during UPDATE, and the third during DELETE.

set_timetravel() allows you to turn time-travel on or off for a table. set_timetravel('mytab', 1) will turn TT ON for table mytab.set_timetravel('mytab', 0) will turn TT OFF for table mytab. In both cases the old status is reported. While TT is off, you can modify the start_date and stop_date columns freely. Note that the on/off status is local to the current database session — fresh sessions will always start out with TT ON for all tables.

get_timetravel() returns the TT state for a table without changing it.

There is an example in timetravel.example.

F.36.3. autoinc — Functions for Autoincrementing Fields

autoinc() is a trigger that stores the next value of a sequence into an integer field. This has some overlap with the built-in “serial column” feature, but it is not the same: autoinc() will override attempts to substitute a different field value during inserts, and optionally it can be used to increment the field during updates, too.

To use, create a BEFORE INSERT (or optionally BEFORE INSERT OR UPDATE) trigger using this function. Specify two trigger arguments: the name of the integer column to be modified, and the name of the sequence object that will supply values. (Actually, you can specify any number of pairs of such names, if you’d like to update more than one autoincrementing column.)

There is an example in autoinc.example.

F.36.4. insert_username — Functions for Tracking Who Changed a Table

insert_username() is a trigger that stores the current user's name into a text field. This can be useful for tracking who last modified a particular row within a table.
Additional Supplied Modules

To use, create a `BEFORE INSERT` and/or `UPDATE` trigger using this function. Specify a single trigger argument: the name of the text column to be modified.

There is an example in `insert_username.example`.

F.36.5. moddatetime — Functions for Tracking Last Modification Time

`moddatetime()` is a trigger that stores the current time into a `timestamp` field. This can be useful for tracking the last modification time of a particular row within a table.

To use, create a `BEFORE UPDATE` trigger using this function. Specify a single trigger argument: the name of the column to be modified. The column must be of type `timestamp` or `timestamp with time zone`.

There is an example in `moddatetime.example`.

F.37. sslinfo

The `sslinfo` module provides information about the SSL certificate that the current client provided when connecting to PostgreSQL. The module is useless (most functions will return NULL) if the current connection does not use SSL.

This extension won't build at all unless the installation was configured with `--with-openssl`.

F.37.1. Functions Provided

`ssl_is_used()` returns boolean

Returns true if current connection to server uses SSL, and false otherwise.

`ssl_version()` returns text

Returns the name of the protocol used for the SSL connection (e.g. TLSv1.0 TLSv1.1, or TLSv1.2).

`ssl_cipher()` returns text

Returns the name of the cipher used for the SSL connection (e.g. DHE-RSA-AES256-SHA).

`ssl_client_cert_present()` returns boolean

Returns true if current client has presented a valid SSL client certificate to the server, and false otherwise. (The server might or might not be configured to require a client certificate.)

`ssl_client_serial()` returns numeric

Returns serial number of current client certificate. The combination of certificate serial number and certificate issuer is guaranteed to uniquely identify a certificate (but not its owner — the owner ought to regularly change their keys, and get new certificates from the issuer).

So, if you run your own CA and allow only certificates from this CA to be accepted by the server, the serial number is the most reliable (albeit not very mnemonic) means to identify a user.

`ssl_client_dn()` returns text

Returns the full subject of the current client certificate, converting character data into the current database encoding. It is assumed that if you use non-ASCII characters in the certificate names,
your database is able to represent these characters, too. If your database uses the SQL_ASCII encoding, non-ASCII characters in the name will be represented as UTF-8 sequences.

The result looks like /CN=Somebody /C=Some country/O=Some organization.

ssl_issuer_dn() returns text

Returns the full issuer name of the current client certificate, converting character data into the current database encoding. Encoding conversions are handled the same as for ssl_client_dn.

The combination of the return value of this function with the certificate serial number uniquely identifies the certificate.

This function is really useful only if you have more than one trusted CA certificate in your server's certificate authority file, or if this CA has issued some intermediate certificate authority certificates.

ssl_client_dn_field(fieldname text) returns text

This function returns the value of the specified field in the certificate subject, or NULL if the field is not present. Field names are string constants that are converted into ASN1 object identifiers using the OpenSSL object database. The following values are acceptable:

- commonName (alias CN)
- surname (alias SN)
- givenName (alias GN)
- countryName (alias C)
- localityName (alias L)
- stateOrProvinceName (alias ST)
- organizationName (alias O)
- organizationUnitName (alias OU)
- title
- description
- initials
- postalCode
- streetAddress
- generationQualifier
- description
- dnQualifier
- x500UniqueIdentifier
- pseudonym
- role
- emailAddress

All of these fields are optional, except commonName. It depends entirely on your CA's policy which of them would be included and which wouldn't. The meaning of these fields, however, is strictly defined by the X.500 and X.509 standards, so you cannot just assign arbitrary meaning to them.

ssl_issuer_field(fieldname text) returns text

Same as ssl_client_dn_field, but for the certificate issuer rather than the certificate subject.

ssl_extension_info() returns setof record

Provide information about extensions of client certificate: extension name, extension value, and if it is a critical extension.
F.37.2. Author

Victor Wagner <vitus@cryptocom.ru>, Cryptocom LTD

Dmitry Voronin <carriingfate92@yandex.ru>

E-Mail of Cryptocom OpenSSL development group: <openssl@cryptocom.ru>

F.38. tablefunc

The tablefunc module includes various functions that return tables (that is, multiple rows). These functions are useful both in their own right and as examples of how to write C functions that return multiple rows.

F.38.1. Functions Provided

Table F.30 shows the functions provided by the tablefunc module.

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>normal_rand(int numvals, float8 mean, float8 stddev)</td>
<td>setof float8</td>
<td>Produces a set of normally distributed random values</td>
</tr>
<tr>
<td>crosstab(text sql)</td>
<td>setof record</td>
<td>Produces a “pivot table” containing row names plus N value columns, where N is determined by the row type specified in the calling query</td>
</tr>
<tr>
<td>crosstabN(text sql)</td>
<td>setof table_crosstab_N</td>
<td>Produces a “pivot table” containing row names plus N value columns. crosstab2, crosstab3, and crosstab4 are predefined, but you can create additional crosstabN functions as described below</td>
</tr>
<tr>
<td>crosstab(text source_sql, text category_sql)</td>
<td>setof record</td>
<td>Produces a “pivot table” with the value columns specified by a second query</td>
</tr>
<tr>
<td>crosstab(text sql, int N)</td>
<td>setof record</td>
<td>Obsolete version of crosstab(text). The parameter N is now ignored, since the number of value columns is always determined by the calling query</td>
</tr>
<tr>
<td>connectby(text relname, text keyid_fld, text parent_keyid_fld [, text orderby_fld ], text start_with, int max_depth [, text branch_delim ])</td>
<td>setof record</td>
<td>Produces a representation of a hierarchical tree structure</td>
</tr>
</tbody>
</table>

F.38.1.1. normal_rand
normal_rand(int numvals, float8 mean, float8 stddev) returns setof float8

normal_rand produces a set of normally distributed random values (Gaussian distribution).

numvals is the number of values to be returned from the function. mean is the mean of the normal distribution of values and stddev is the standard deviation of the normal distribution of values.

For example, this call requests 1000 values with a mean of 5 and a standard deviation of 3:

test=# SELECT * FROM normal_rand(1000, 5, 3);

<table>
<thead>
<tr>
<th>normal_rand</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.56556322244898</td>
</tr>
<tr>
<td>9.10040991424657</td>
</tr>
<tr>
<td>5.36957140345079</td>
</tr>
<tr>
<td>-0.369151492880995</td>
</tr>
<tr>
<td>0.283600703686639</td>
</tr>
<tr>
<td>.</td>
</tr>
<tr>
<td>.</td>
</tr>
<tr>
<td>4.82992125404908</td>
</tr>
<tr>
<td>9.71308014517282</td>
</tr>
<tr>
<td>2.49639286969028</td>
</tr>
</tbody>
</table>

(1000 rows)

F.38.1.2. crosstab (text)

crosstab(text sql)
crosstab(text sql, int N)

The crosstab function is used to produce “pivot” displays, wherein data is listed across the page rather than down. For example, we might have data like

```
row1  val11
row1  val12
row1  val13
...
row2  val21
row2  val22
row2  val23
...
```

which we wish to display like

```
row1  val11  val12  val13  ...
row2  val21  val22  val23  ...
...
```

The crosstab function takes a text parameter that is a SQL query producing raw data formatted in the first way, and produces a table formatted in the second way.

The sql parameter is a SQL statement that produces the source set of data. This statement must return one row_name column, one category column, and one value column. N is an obsolete parameter, ignored if supplied (formerly this had to match the number of output value columns, but now that is determined by the calling query).
For example, the provided query might produce a set something like:

<table>
<thead>
<tr>
<th>row_name</th>
<th>cat</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>row1</td>
<td>cat1</td>
<td>val1</td>
</tr>
<tr>
<td>row1</td>
<td>cat2</td>
<td>val2</td>
</tr>
<tr>
<td>row1</td>
<td>cat3</td>
<td>val3</td>
</tr>
<tr>
<td>row1</td>
<td>cat4</td>
<td>val4</td>
</tr>
<tr>
<td>row2</td>
<td>cat1</td>
<td>val5</td>
</tr>
<tr>
<td>row2</td>
<td>cat2</td>
<td>val6</td>
</tr>
<tr>
<td>row2</td>
<td>cat3</td>
<td>val7</td>
</tr>
<tr>
<td>row2</td>
<td>cat4</td>
<td>val8</td>
</tr>
</tbody>
</table>

The \texttt{crosstab} function is declared to return \texttt{setof record}, so the actual names and types of the output columns must be defined in the \texttt{FROM} clause of the calling \texttt{SELECT} statement, for example:

\begin{verbatim}
SELECT * FROM crosstab('...') AS ct(row_name text, category_1 text, category_2 text);
\end{verbatim}

This example produces a set something like:

<table>
<thead>
<tr>
<th>row_name</th>
<th>category_1</th>
<th>category_2</th>
</tr>
</thead>
<tbody>
<tr>
<td>row1</td>
<td>val1</td>
<td>val2</td>
</tr>
<tr>
<td>row2</td>
<td>val5</td>
<td>val6</td>
</tr>
</tbody>
</table>

The \texttt{FROM} clause must define the output as one \texttt{row_name} column (of the same data type as the first result column of the SQL query) followed by \texttt{N value} columns (all of the same data type as the third result column of the SQL query). You can set up as many output value columns as you wish. The names of the output columns are up to you.

The \texttt{crosstab} function produces one output row for each consecutive group of input rows with the same \texttt{row_name} value. It fills the output value columns, left to right, with the value fields from these rows. If there are fewer rows in a group than there are output value columns, the extra output columns are filled with nulls; if there are more rows, the extra input rows are skipped.

In practice the SQL query should always specify \texttt{ORDER BY 1,2} to ensure that the input rows are properly ordered, that is, values with the same \texttt{row_name} are brought together and correctly ordered within the row. Notice that \texttt{crosstab} itself does not pay any attention to the second column of the query result; it’s just there to be ordered by, to control the order in which the third-column values appear across the page.

Here is a complete example:

\begin{verbatim}
CREATE TABLE ct(id SERIAL, rowid TEXT, attribute TEXT, value TEXT);
INSERT INTO ct(rowid, attribute, value)
VALUES('test1','att1','val1');
INSERT INTO ct(rowid, attribute, value)
VALUES('test1','att2','val2');
INSERT INTO ct(rowid, attribute, value)
VALUES('test1','att3','val3');
INSERT INTO ct(rowid, attribute, value)
VALUES('test1','att4','val4');
INSERT INTO ct(rowid, attribute, value)
VALUES('test2','att1','val5');
\end{verbatim}
INSERT INTO ct(rowid, attribute, value)
VALUES('test2','att2','val6');
INSERT INTO ct(rowid, attribute, value)
VALUES('test2','att3','val7');
INSERT INTO ct(rowid, attribute, value)
VALUES('test2','att4','val8');

SELECT *
FROM crosstab(
    'select rowid, attribute, value
    from ct
    where attribute = ''att2'' or attribute = ''att3''
    order by 1,2')
AS ct(row_name text, category_1 text, category_2 text, category_3 text);

<table>
<thead>
<tr>
<th>row_name</th>
<th>category_1</th>
<th>category_2</th>
<th>category_3</th>
</tr>
</thead>
<tbody>
<tr>
<td>test1</td>
<td>val2</td>
<td>val3</td>
<td></td>
</tr>
<tr>
<td>test2</td>
<td>val6</td>
<td>val7</td>
<td></td>
</tr>
</tbody>
</table>

(2 rows)

You can avoid always having to write out a FROM clause to define the output columns, by setting up a custom crosstab function that has the desired output row type wired into its definition. This is described in the next section. Another possibility is to embed the required FROM clause in a view definition.

Note

See also the \crosstabview command in psql, which provides functionality similar to crosstab().

F.38.1.3. crosstabN(text)

crosstabN(text sql)

The crosstabN functions are examples of how to set up custom wrappers for the general crosstab function, so that you need not write out column names and types in the calling SELECT query. The tablefunc module includes crosstab2, crosstab3, and crosstab4, whose output row types are defined as

CREATE TYPE tablefunc_crosstab_N AS (
    row_name TEXT,
    category_1 TEXT,
    category_2 TEXT,
    .
    .
    category_N TEXT
);

Thus, these functions can be used directly when the input query produces row_name and value columns of type text, and you want 2, 3, or 4 output values columns. In all other ways they behave exactly as described above for the general crosstab function.

For instance, the example given in the previous section would also work as
SELECT * 
FROM crosstab3(
    'select rowid, attribute, value
    from ct
    where attribute = ''att2'' or attribute = ''att3''
    order by 1,2');

These functions are provided mostly for illustration purposes. You can create your own return types and functions based on the underlying crosstab() function. There are two ways to do it:

- Create a composite type describing the desired output columns, similar to the examples in contrib/tablefunc/tablefunc--1.0.sql. Then define a unique function name accepting one text parameter and returning setof your_type_name, but linking to the same underlying crosstab C function. For example, if your source data produces row names that are text, and values that are float8, and you want 5 value columns:

  CREATE TYPE my_crosstab_float8_5_cols AS (my_row_name text,
                                           my_category_1 float8,
                                           my_category_2 float8,
                                           my_category_3 float8,
                                           my_category_4 float8,
                                           my_category_5 float8);

  CREATE OR REPLACE FUNCTION crosstab_float8_5_cols(text) RETURNS setof my_crosstab_float8_5_cols
                  AS '$libdir/tablefunc','crosstab' LANGUAGE C STABLE STRICT;

- Use OUT parameters to define the return type implicitly. The same example could also be done this way:

  CREATE OR REPLACE FUNCTION crosstab_float8_5_cols(
      IN text,
      OUT my_row_name text,
      OUT my_category_1 float8,
      OUT my_category_2 float8,
      OUT my_category_3 float8,
      OUT my_category_4 float8,
      OUT my_category_5 float8)
  RETURNS setof record
  AS '$libdir/tablefunc','crosstab' LANGUAGE C STABLE STRICT;

**F.38.1.4. crosstab(text, text)**

crosstab(text source_sql, text category_sql)

The main limitation of the single-parameter form of crosstab is that it treats all values in a group alike, inserting each value into the first available column. If you want the value columns to correspond to specific categories of data, and some groups might not have data for some of the categories, that doesn't work well. The two-parameter form of crosstab handles this case by providing an explicit list of the categories corresponding to the output columns.

source_sql is a SQL statement that produces the source set of data. This statement must return one row_name column, one category column, and one value column. It may also have one or more
“extra” columns. The row_name column must be first. The category and value columns must be
the last two columns, in that order. Any columns between row_name and category are treated as
“extra”. The “extra” columns are expected to be the same for all rows with the same row_name value.

For example, source_sql might produce a set something like:

```
SELECT row_name, extra_col, cat, value FROM foo ORDER BY 1;
```

```
row_name    extra_col   cat    value
----------+------------+-----+---------
row1         extra1    cat1    val1
row1         extra1    cat2    val2
row1         extra1    cat4    val4
row2         extra2    cat1    val5
row2         extra2    cat2    val6
row2         extra2    cat3    val7
row2         extra2    cat4    val8
```

category_sql is a SQL statement that produces the set of categories. This statement must return
only one column. It must produce at least one row, or an error will be generated. Also, it must not
produce duplicate values, or an error will be generated. category_sql might be something like:

```
SELECT DISTINCT cat FROM foo ORDER BY 1;
```

```
cat
------
cat1
cat2
cat3
cat4
```

The crosstab function is declared to return setof record, so the actual names and types of the
output columns must be defined in the FROM clause of the calling SELECT statement, for example:

```
SELECT * FROM crosstab('...', '...')
    AS ct(row_name text, extra text, cat1 text, cat2 text, cat3 text, cat4 text);
```

This will produce a result something like:

```
<== value columns ==> 
row_name    extra    cat1   cat2   cat3   cat4
----------+--------+------+------+-------+
row1        extra1  val1   val2          val4
row2        extra2  val5   val6   val7   val8
```

The FROM clause must define the proper number of output columns of the proper data types. If there
are N columns in the source_sql query's result, the first N-2 of them must match up with the first
N-2 output columns. The remaining output columns must have the type of the last column of the
source_sql query's result, and there must be exactly as many of them as there are rows in the
category_sql query's result.

The crosstab function produces one output row for each consecutive group of input rows with the
same row_name value. The output row_name column, plus any “extra” columns, are copied from
the first row of the group. The output value columns are filled with the value fields from rows
having matching category values. If a row's category does not match any output of the cate-
**Additional Supplied Modules**

*gory_sql* query, its value is ignored. Output columns whose matching category is not present in any input row of the group are filled with nulls.

In practice the *source_sql* query should always specify ORDER BY 1 to ensure that values with the same row_name are brought together. However, ordering of the categories within a group is not important. Also, it is essential to be sure that the order of the *category_sql* query's output matches the specified output column order.

Here are two complete examples:

```sql
create table sales(year int, month int, qty int);
insert into sales values(2007, 1, 1000);
insert into sales values(2007, 2, 1500);
insert into sales values(2007, 7, 500);
insert into sales values(2007, 11, 1500);
insert into sales values(2007, 12, 2000);
insert into sales values(2008, 1, 1000);

select * from crosstab("select year, month, qty from sales order by 1",
                        "select m from generate_series(1,12) m")
  as (year int, "Jan" int, "Feb" int, "Mar" int, "Apr" int, "May" int, "Jun" int, "Jul" int, "Aug" int, "Sep" int, "Oct" int, "Nov" int, "Dec" int)
);  

year | Jan  | Feb  | Mar | Apr | May | Jun | Jul | Aug | Sep | Oct  
|-----|------|------|-----|-----|-----|-----|-----|-----|-----|-----  
2007 | 1000 | 1500 |     |     |     |     | 500 |     |     |       
| 1500 | 2000  
2008 | 1000 |      |     |     |     |     |     |     |     |       
|      |       
(2 rows)

CREATE TABLE cth(rowid text, rowdt timestamp, attribute text, val text);
INSERT INTO cth VALUES('test1','01 March 2003','temperature','42');
INSERT INTO cth VALUES('test1','01 March 2003','test_result','PASS');
INSERT INTO cth VALUES('test2','02 March 2003','temperature','53');
INSERT INTO cth VALUES('test2','02 March 2003','test_startdate','01 March 2003');
```

3441
INSERT INTO cth VALUES('test2','02 March 2003','volts','3.1234');

SELECT * FROM crosstab
(
    'SELECT rowid, rowdt, attribute, val FROM cth ORDER BY 1',
    'SELECT DISTINCT attribute FROM cth ORDER BY 1'
)
AS
(
    rowid text,
    rowdt timestamp,
    temperature int4,
    test_result text,
    test_startdate timestamp,
    volts float8
);

| rowid |          rowdt           | temperature | test_result |
|-------+--------------------------+-------------+-------------|
|       | test_startdate | volts       |
|-------+--------------------------+-------------+-------------|
| test1 | Sat Mar 01 00:00:00 2003 | 42          | PASS        |
|       | 2.6987                  |             |
| test2 | Sun Mar 02 00:00:00 2003 | 53          | FAIL        | Sat Mar 01 00:00:00 2003 | 3.1234

(2 rows)

You can create predefined functions to avoid having to write out the result column names and types in each query. See the examples in the previous section. The underlying C function for this form of crosstab is named crosstab_hash.

F.38.1.5. connectby

connectby(text relname, text keyid_fld, text parent_keyid_fld
  [, text orderby_fld ], text start_with, int max_depth
  [, text branch_delim ])

The connectby function produces a display of hierarchical data that is stored in a table. The table must have a key field that uniquely identifies rows, and a parent-key field that references the parent (if any) of each row. connectby can display the sub-tree descending from any row.

Table F.31 explains the parameters.

Table F.31. connectby Parameters

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relname</td>
<td>Name of the source relation</td>
</tr>
<tr>
<td>keyid_fld</td>
<td>Name of the key field</td>
</tr>
<tr>
<td>parent_keyid_fld</td>
<td>Name of the parent-key field</td>
</tr>
<tr>
<td>orderby_fld</td>
<td>Name of the field to order siblings by (optional)</td>
</tr>
<tr>
<td>start_with</td>
<td>Key value of the row to start at</td>
</tr>
<tr>
<td>max_depth</td>
<td>Maximum depth to descend to, or zero for unlimited depth</td>
</tr>
<tr>
<td>branch_delim</td>
<td>String to separate keys with in branch output (optional)</td>
</tr>
</tbody>
</table>
The key and parent-key fields can be any data type, but they must be the same type. Note that the `start_with` value must be entered as a text string, regardless of the type of the key field.

The `connectby` function is declared to return `setof record`, so the actual names and types of the output columns must be defined in the `FROM` clause of the calling `SELECT` statement, for example:

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'pos', 'row2', 0, '~')
  AS t(keyid text, parent_keyid text, level int, branch text, pos int);
```

The first two output columns are used for the current row's key and its parent row's key; they must match the type of the table's key field. The third output column is the depth in the tree and must be of type `integer`. If a `branch_delim` parameter was given, the next output column is the branch display and must be of type `text`. Finally, if an `orderby_fld` parameter was given, the last output column is a serial number, and must be of type `integer`.

The “branch” output column shows the path of keys taken to reach the current row. The keys are separated by the specified `branch_delim` string. If no branch display is wanted, omit both the `branch_delim` parameter and the branch column in the output column list.

If the ordering of siblings of the same parent is important, include the `orderby_fld` parameter to specify which field to order siblings by. This field can be of any sortable data type. The output column list must include a final integer serial-number column, if and only if `orderby_fld` is specified.

The parameters representing table and field names are copied as-is into the SQL queries that `connectby` generates internally. Therefore, include double quotes if the names are mixed-case or contain special characters. You may also need to schema-qualify the table name.

In large tables, performance will be poor unless there is an index on the parent-key field.

It is important that the `branch_delim` string not appear in any key values, else `connectby` may incorrectly report an infinite-recursion error. Note that if `branch_delim` is not provided, a default value of `~` is used for recursion detection purposes.

Here is an example:

```sql
CREATE TABLE connectby_tree(keyid text, parent_keyid text, pos int);

INSERT INTO connectby_tree VALUES('row1',NULL, 0);
INSERT INTO connectby_tree VALUES('row2','row1', 0);
INSERT INTO connectby_tree VALUES('row3','row1', 0);
INSERT INTO connectby_tree VALUES('row4','row2', 1);
INSERT INTO connectby_tree VALUES('row5','row2', 0);
INSERT INTO connectby_tree VALUES('row6','row4', 0);
INSERT INTO connectby_tree VALUES('row7','row3', 0);
INSERT INTO connectby_tree VALUES('row8','row6', 0);
INSERT INTO connectby_tree VALUES('row9','row5', 0);

-- with branch, without orderby_fld (order of results is not guaranteed)
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'row2', 0, '~')
  AS t(keyid text, parent_keyid text, level int, branch text);
keyid | parent_keyid | level | branch
-------+--------------+-------+---------------------
row2  |              |     0 | row2
row4  | row2         |     1 | row2~row4
```

3443
### Additional Supplied Modules

<table>
<thead>
<tr>
<th>row6</th>
<th>row4</th>
<th>2</th>
<th>row2-row4-row6</th>
</tr>
</thead>
<tbody>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
<td>row2-row4-row6-row8</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
<td>row2-row5</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
<td>row2-row5-row9</td>
</tr>
</tbody>
</table>

(6 rows)

-- without branch, without orderby_fld (order of results is not guaranteed)

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'row2', 0)
AS t(keyid text, parent_keyid text, level int);
```

<table>
<thead>
<tr>
<th>keyid</th>
<th>parent_keyid</th>
<th>level</th>
</tr>
</thead>
<tbody>
<tr>
<td>row2</td>
<td></td>
<td>0</td>
</tr>
<tr>
<td>row4</td>
<td>row2</td>
<td>1</td>
</tr>
<tr>
<td>row6</td>
<td>row4</td>
<td>2</td>
</tr>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
</tr>
</tbody>
</table>

(6 rows)

-- with branch, with orderby_fld (notice that row5 comes before row4)

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'pos', 'row2', 0, '~')
AS t(keyid text, parent_keyid text, level int, branch text, pos int);
```

<table>
<thead>
<tr>
<th>keyid</th>
<th>parent_keyid</th>
<th>level</th>
<th>branch</th>
<th>pos</th>
</tr>
</thead>
<tbody>
<tr>
<td>row2</td>
<td></td>
<td>0</td>
<td>row2</td>
<td>1</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
<td>row2~row5</td>
<td>2</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
<td>row2<del>row5</del>row9</td>
<td>3</td>
</tr>
<tr>
<td>row4</td>
<td>row2</td>
<td>1</td>
<td>row2~row4</td>
<td>4</td>
</tr>
<tr>
<td>row6</td>
<td>row4</td>
<td>2</td>
<td>row2<del>row4</del>row6</td>
<td>5</td>
</tr>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
<td>row2<del>row4</del>row6~row8</td>
<td>6</td>
</tr>
</tbody>
</table>

(6 rows)

-- without branch, with orderby_fld (notice that row5 comes before row4)

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'pos', 'row2', 0)
AS t(keyid text, parent_keyid text, level int, pos int);
```

<table>
<thead>
<tr>
<th>keyid</th>
<th>parent_keyid</th>
<th>level</th>
<th>pos</th>
</tr>
</thead>
<tbody>
<tr>
<td>row2</td>
<td></td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
<td>2</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td>row4</td>
<td>row2</td>
<td>1</td>
<td>4</td>
</tr>
<tr>
<td>row6</td>
<td>row4</td>
<td>2</td>
<td>5</td>
</tr>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
<td>6</td>
</tr>
</tbody>
</table>

(6 rows)

---

**F.38.2. Author**

Joe Conway

**F.39. tcn**
The `tcn` module provides a trigger function that notifies listeners of changes to any table on which it is attached. It must be used as an `AFTER` trigger `FOR EACH ROW`.

Only one parameter may be supplied to the function in a `CREATE TRIGGER` statement, and that is optional. If supplied it will be used for the channel name for the notifications. If omitted `tcn` will be used for the channel name.

The payload of the notifications consists of the table name, a letter to indicate which type of operation was performed, and column name/value pairs for primary key columns. Each part is separated from the next by a comma. For ease of parsing using regular expressions, table and column names are always wrapped in double quotes, and data values are always wrapped in single quotes. Embedded quotes are doubled.

A brief example of using the extension follows.

```sql
CREATE TABLE test=# create table tcndata
test=# (test(# a int not null,
test(# b date not null,
test(# c text,
test(# primary key (a, b)
test(# );
test=# create trigger tcndata_tcn_trigger
test=# after insert or update or delete on tcndata
  test=# for each row execute function
  triggered_change_notification();
test=# listen tcn;
LISTEN
test=# insert into tcndata values (1, date '2012-12-22', 'one'),
test=# (1, date '2012-12-23',
  'another'),
test=# (2, date '2012-12-23', 'two');
INSERT 0 3
Asynchronous notification "tcn" with payload
  ""tcndata",I,"a"='1','b'='2012-12-22'" received from server
  process with PID 22770.
Asynchronous notification "tcn" with payload
  ""tcndata",I,"a"='1','b'='2012-12-23'" received from server
  process with PID 22770.
Asynchronous notification "tcn" with payload
  ""tcndata",I,"a"='2','b'='2012-12-23'" received from server
  process with PID 22770.
test=# update tcndata set c = 'uno' where a = 1;
UPDATE 2
Asynchronous notification "tcn" with payload
  ""tcndata",U,"a"='1','b'='2012-12-22'" received from server
  process with PID 22770.
Asynchronous notification "tcn" with payload
  ""tcndata",U,"a"='1','b'='2012-12-23'" received from server
  process with PID 22770.
test=# delete from tcndata where a = 1 and b = date '2012-12-22';
DELETE 1
Asynchronous notification "tcn" with payload
  ""tcndata",D,"a"='1','b'='2012-12-22'" received from server
  process with PID 22770.
```

3445
test_decoding is an example of a logical decoding output plugin. It doesn't do anything especially useful, but can serve as a starting point for developing your own output plugin.

test_decoding receives WAL through the logical decoding mechanism and decodes it into text representations of the operations performed.

Typical output from this plugin, used over the SQL logical decoding interface, might be:

```
postgres=# SELECT * FROM pg_logical_slot_get_changes('test_slot',
NULL, NULL, 'include-xids', '0');
  lsn | xid |                       data
---------------------
0/16D30F8 | 691 | BEGIN
0/16D32A0 | 691 | table public.data: INSERT: id[int4]:2
data[text]:'arg'
0/16D32A0 | 691 | table public.data: INSERT: id[int4]:3
data[text]:'demo'
0/16D32A0 | 691 | COMMIT
0/16D32D8 | 692 | BEGIN
0/16D3398 | 692 | table public.data: DELETE: id[int4]:2
0/16D3398 | 692 | table public.data: DELETE: id[int4]:3
0/16D3398 | 692 | COMMIT
```

**F.41. tsm_system_rows**

The tsm_system_rows module provides the table sampling method SYSTEM_ROWS, which can be used in the TABLESAMPLE clause of a SELECT command.

This table sampling method accepts a single integer argument that is the maximum number of rows to read. The resulting sample will always contain exactly that many rows, unless the table does not contain enough rows, in which case the whole table is selected.

Like the built-in SYSTEM sampling method, SYSTEM_ROWS performs block-level sampling, so that the sample is not completely random but may be subject to clustering effects, especially if only a small number of rows are requested.

SYSTEM_ROWS does not support the REPEATABLE clause.

**F.41.1. Examples**

Here is an example of selecting a sample of a table with SYSTEM_ROWS. First install the extension:

```
CREATE EXTENSION tsm_system_rows;
```

Then you can use it in a SELECT command, for instance:

```
SELECT * FROM my_table TABLESAMPLE SYSTEM_ROWS(100);
```

This command will return a sample of 100 rows from the table my_table (unless the table does not have 100 visible rows, in which case all its rows are returned).

**F.42. tsm_system_time**
The `tsm_system_time` module provides the table sampling method `SYSTEM_TIME`, which can be used in the `TABLESAMPLE` clause of a `SELECT` command.

This table sampling method accepts a single floating-point argument that is the maximum number of milliseconds to spend reading the table. This gives you direct control over how long the query takes, at the price that the size of the sample becomes hard to predict. The resulting sample will contain as many rows as could be read in the specified time, unless the whole table has been read first.

Like the built-in `SYSTEM` sampling method, `SYSTEM_TIME` performs block-level sampling, so that the sample is not completely random but may be subject to clustering effects, especially if only a small number of rows are selected.

`SYSTEM_TIME` does not support the `REPEATABLE` clause.

**F.42.1. Examples**

Here is an example of selecting a sample of a table with `SYSTEM_TIME`. First install the extension:

```sql
CREATE EXTENSION tsm_system_time;
```

Then you can use it in a `SELECT` command, for instance:

```sql
SELECT * FROM my_table TABLESAMPLE SYSTEM_TIME(1000);
```

This command will return as large a sample of `my_table` as it can read in 1 second (1000 milliseconds). Of course, if the whole table can be read in under 1 second, all its rows will be returned.

**F.43. unaccent**

`unaccent` is a text search dictionary that removes accents (diacritic signs) from lexemes. It's a filtering dictionary, which means its output is always passed to the next dictionary (if any), unlike the normal behavior of dictionaries. This allows accent-insensitive processing for full text search.

The current implementation of `unaccent` cannot be used as a normalizing dictionary for the thesaurus dictionary.

**F.43.1. Configuration**

An `unaccent` dictionary accepts the following options:

- `RULES` is the base name of the file containing the list of translation rules. This file must be stored in `$SHAREDIR/tsearch_data/` (where `$SHAREDIR` means the PostgreSQL installation's shared-data directory). Its name must end in `.rules` (which is not to be included in the `RULES` parameter).

The rules file has the following format:

- Each line represents one translation rule, consisting of a character with accent followed by a character without accent. The first is translated into the second. For example,

```
À A
Á A
Ã A
Å A
Æ AE
```
The two characters must be separated by whitespace, and any leading or trailing whitespace on a line is ignored.

- Alternatively, if only one character is given on a line, instances of that character are deleted; this is useful in languages where accents are represented by separate characters.
- Actually, each “character” can be any string not containing whitespace, so unaccent dictionaries could be used for other sorts of substring substitutions besides diacritic removal.
- As with other PostgreSQL text search configuration files, the rules file must be stored in UTF-8 encoding. The data is automatically translated into the current database's encoding when loaded. Any lines containing untranslatable characters are silently ignored, so that rules files can contain rules that are not applicable in the current encoding.

A more complete example, which is directly useful for most European languages, can be found in unaccent.rules, which is installed in $SHAREDIR/tsearch_data/ when the unaccent module is installed. This rules file translates characters with accents to the same characters without accents, and it also expands ligatures into the equivalent series of simple characters (for example, Æ to AE).

\section*{F.43.2. Usage}

Installing the unaccent extension creates a text search template unaccent and a dictionary unaccent based on it. The unaccent dictionary has the default parameter setting RULES='unaccent', which makes it immediately usable with the standard unaccent.rules file. If you wish, you can alter the parameter, for example

```
mydb=# ALTER TEXT SEARCH DICTIONARY unaccent (RULES='my_rules');
```
or create new dictionaries based on the template.

To test the dictionary, you can try:

```
mydb=# select ts_lexize('unaccent','Hôtel');
  ts_lexize
     -----------------------
     {Hotel}
(1 row)
```

Here is an example showing how to insert the unaccent dictionary into a text search configuration:

```
mydb=# CREATE TEXT SEARCH CONFIGURATION fr ( COPY = french );
mydb=# ALTER TEXT SEARCH CONFIGURATION fr
  ALTER MAPPING FOR hword, hword_part, word
    WITH unaccent, french_stem;
mydb=# select to_tsvector('fr','Hôtels de la Mer');
  to_tsvector
     ----------------------
     'hotel':1 'mer':4
(1 row)
```

```
mydb=# select to_tsvector('fr','Hôtel de la Mer') @@
     to_tsquery('fr','Hotels');
?column?
  ---------
  t
(1 row)
```
Additional Supplied Modules

mydb=# select ts_headline('fr','Hôtel de la Mer',to_tsquery('fr','Hotels'));

<table>
<thead>
<tr>
<th>ts_headline</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;b&gt;Hôtel&lt;/b&gt; de la Mer</td>
</tr>
</tbody>
</table>

(1 row)

**F.43.3. Functions**

The `unaccent()` function removes accents (diacritic signs) from a given string. Basically, it’s a wrapper around `unaccent`-type dictionaries, but it can be used outside normal text search contexts.

\[
\text{unaccent([dictionary regdictionary, ] string text) returns text}
\]

If the `dictionary` argument is omitted, the text search dictionary named `unaccent` and appearing in the same schema as the `unaccent()` function itself is used.

For example:

```sql
SELECT unaccent('unaccent', 'Hôtel');
SELECT unaccent('Hôtel');
```

**F.44. uuid-ossp**

The `uuid-ossp` module provides functions to generate universally unique identifiers (UUIDs) using one of several standard algorithms. There are also functions to produce certain special UUID constants.

**F.44.1. uuid-ossp Functions**

Table F.32 shows the functions available to generate UUIDs. The relevant standards ITU-T Rec. X.667, ISO/IEC 9834-8:2005, and RFC 4122 specify four algorithms for generating UUIDs, identified by the version numbers 1, 3, 4, and 5. (There is no version 2 algorithm.) Each of these algorithms could be suitable for a different set of applications.

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uuid_generate_v1()</td>
<td>This function generates a version 1 UUID. This involves the MAC address of the computer and a time stamp. Note that UUIDs of this kind reveal the identity of the computer that created the identifier and the time at which it did so, which might make it unsuitable for certain security-sensitive applications.</td>
</tr>
<tr>
<td>uuid_generate_v1mc()</td>
<td>This function generates a version 1 UUID but uses a random multicast MAC address instead of the real MAC address of the computer.</td>
</tr>
<tr>
<td>uuid_generate_v3(namespace uuid, name text)</td>
<td>This function generates a version 3 UUID in the given namespace using the specified input name. The namespace should be one of the special constants produced by the <code>uuid_ns_*()</code> functions shown in Table F.33. (It could be any UUID in theory.) The name is an identifier in the selected namespace.</td>
</tr>
</tbody>
</table>
Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
</table>
| uuid_generate_v3() | For example:  
  ```sql
  SELECT uuid_generate_v3(uuid_ns_url(), 'http://www.postgresql.org');
  ```  
The name parameter will be MD5-hashed, so the cleartext cannot be derived from the generated UUID. The generation of UUIDs by this method has no random or environment-dependent element and is therefore reproducible. |
| uuid_generate_v4() | This function generates a version 4 UUID, which is derived entirely from random numbers. |
| uuid_generate_v5(namespace uuid, name text) | This function generates a version 5 UUID, which works like a version 3 UUID except that SHA-1 is used as a hashing method. Version 5 should be preferred over version 3 because SHA-1 is thought to be more secure than MD5. |

Table F.33. Functions Returning UUID Constants

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uuid_nil()</td>
<td>A “nil” UUID constant, which does not occur as a real UUID.</td>
</tr>
<tr>
<td>uuid_ns_dns()</td>
<td>Constant designating the DNS namespace for UUIDs.</td>
</tr>
<tr>
<td>uuid_ns_url()</td>
<td>Constant designating the URL namespace for UUIDs.</td>
</tr>
<tr>
<td>uuid_ns_oid()</td>
<td>Constant designating the ISO object identifier (OID) namespace for UUIDs. (This pertains to ASN.1 OIDs, which are unrelated to the OIDs used in PostgreSQL.)</td>
</tr>
<tr>
<td>uuid_ns_x500()</td>
<td>Constant designating the X.500 distinguished name (DN) namespace for UUIDs.</td>
</tr>
</tbody>
</table>

F.44.2. Building uuid-ossp

Historically this module depended on the OSSP UUID library, which accounts for the module’s name. While the OSSP UUID library can still be found at http://www.ossp.org/pkg/lib/uuid/, it is not well maintained, and is becoming increasingly difficult to port to newer platforms. uuid-ossp can now be built without the OSSP library on some platforms. On FreeBSD, NetBSD, and some other BSD-derived platforms, suitable UUID creation functions are included in the core libc library. On Linux, macOS, and some other platforms, suitable functions are provided in the libuuid library, which originally came from the e2fsprogs project (though on modern Linux it is considered part of util-linux-ng). When invoking configure, specify --with-uuid=bsd to use the BSD functions, or --with-uuid=e2fs to use e2fsprogs' libuuid, or --with-uuid=ossp to use the OSSP UUID library. More than one of these libraries might be available on a particular machine, so configure does not automatically choose one.

Note

If you only need randomly-generated (version 4) UUIDs, consider using the gen_random_uuid() function from the pgcrypto module instead.
F.44.3. Author

Peter Eisentraut <peter_e@gmx.net>

F.45. xml2

The xml2 module provides XPath querying and XSLT functionality.

F.45.1. Deprecation Notice

From PostgreSQL 8.3 on, there is XML-related functionality based on the SQL/XML standard in the core server. That functionality covers XML syntax checking and XPath queries, which is what this module does, and more, but the API is not at all compatible. It is planned that this module will be removed in a future version of PostgreSQL in favor of the newer standard API, so you are encouraged to try converting your applications. If you find that some of the functionality of this module is not available in an adequate form with the newer API, please explain your issue to <pgsql-hackers@postgresql.org> so that the deficiency can be addressed.

F.45.2. Description of Functions

Table F.34 shows the functions provided by this module. These functions provide straightforward XML parsing and XPath queries. All arguments are of type text, so for brevity that is not shown.

Table F.34. Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xml_valid(document)</td>
<td>bool</td>
<td>This parses the document text in its parameter and returns true if the document is well-formed XML. (Note: this is an alias for the standard PostgreSQL function xml_is_well_formed(). The name xml_valid() is technically incorrect since validity and well-formedness have different meanings in XML.)</td>
</tr>
<tr>
<td>xpath_string(document, query)</td>
<td>text</td>
<td>These functions evaluate the XPath query on the supplied document, and cast the result to the specified type.</td>
</tr>
<tr>
<td>xpath_number(document, query)</td>
<td>float4</td>
<td></td>
</tr>
<tr>
<td>xpath_bool(document, query)</td>
<td>bool</td>
<td></td>
</tr>
<tr>
<td>xpath_nodeset(document, query, toptag, itemtag)</td>
<td>text</td>
<td>This evaluates query on document and wraps the result in XML tags. If the result is multi-valued, the output will look like: &lt;toptag&gt;&lt;itemtag&gt;Value 1 which could be an XML fragment&lt;/itemtag&gt;</td>
</tr>
<tr>
<td>Function</td>
<td>Returns</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>---------</td>
<td>-------------</td>
</tr>
<tr>
<td>xpath_nodeset(document, query)</td>
<td>text</td>
<td>Like xpath_nodeset(document, query, toptag, itemtag) but result omits both tags.</td>
</tr>
<tr>
<td>xpath_nodeset(document, query, itemtag)</td>
<td>text</td>
<td>Like xpath_nodeset(document, query, toptag, itemtag) but result omits toptag.</td>
</tr>
<tr>
<td>xpath_list(document, query, separator)</td>
<td>text</td>
<td>This function returns multiple values separated by the specified separator, for example Value 1, Value 2, Value 3 if separator is ,.</td>
</tr>
<tr>
<td>xpath_list(document, query)</td>
<td>text</td>
<td>This is a wrapper for the above function that uses , as the separator.</td>
</tr>
</tbody>
</table>

**F.45.3. xpath_table**

xpath_table(text key, text document, text relation, text xpaths, text criteria) returns setof record

xpath_table is a table function that evaluates a set of XPath queries on each of a set of documents and returns the results as a table. The primary key field from the original document table is returned as the first column of the result so that the result set can readily be used in joins. The parameters are described in Table F.35.

**Table F.35. xpath_table Parameters**

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>key</td>
<td>the name of the “key” field — this is just a field to be used as the first column of the output table, i.e., it identifies the record from which each output row came (see note below about multiple values)</td>
</tr>
<tr>
<td>document</td>
<td>the name of the field containing the XML document</td>
</tr>
<tr>
<td>relation</td>
<td>the name of the table or view containing the documents</td>
</tr>
<tr>
<td>xpaths</td>
<td>one or more XPath expressions, separated by</td>
</tr>
<tr>
<td>criteria</td>
<td>the contents of the WHERE clause. This cannot be omitted, so use true or 1=1 if you want to process all the rows in the relation</td>
</tr>
</tbody>
</table>

These parameters (except the XPath strings) are just substituted into a plain SQL SELECT statement, so you have some flexibility — the statement is
Additional Supplied Modules

SELECT <key>, <document> FROM <relation> WHERE <criteria>

so those parameters can be *anything* valid in those particular locations. The result from this SELECT
needs to return exactly two columns (which it will unless you try to list multiple fields for key or
document). Beware that this simplistic approach requires that you validate any user-supplied values
avoid SQL injection attacks.

The function has to be used in a FROM expression, with an AS clause to specify the output columns;
for example

SELECT * FROM
xpath_table('article_id',
'article_xml',
'articles',
'/article/author|/article/pages|/article/title',
'date_entered > ''2003-01-01'' ')
AS t(article_id integer, author text, page_count integer, title
text);

The AS clause defines the names and types of the columns in the output table. The first is the “key” field
and the rest correspond to the XPath queries. If there are more XPath queries than result columns, the
extra queries will be ignored. If there are more result columns than XPath queries, the extra columns
will be NULL.

Notice that this example defines the page_count result column as an integer. The function deals
internally with string representations, so when you say you want an integer in the output, it will take
the string representation of the XPath result and use PostgreSQL input functions to transform it into an
integer (or whatever type the AS clause requests). An error will result if it can’t do this — for example
if the result is empty — so you may wish to just stick to text as the column type if you think your
data has any problems.

The calling SELECT statement doesn’t necessarily have to be just SELECT * — it can reference the
output columns by name or join them to other tables. The function produces a virtual table with which
you can perform any operation you wish (e.g. aggregation, joining, sorting etc). So we could also have:

SELECT t.title, p.fullname, p.email
FROM xpath_table('article_id', 'article_xml', 'articles',
'/article/title|/article/author/@id',
'xpath_string(article_xml,''/article/@date'') >
'2003-03-20' ')
AS t(article_id integer, title text, author_id integer),
thlPeopleInfo AS p
WHERE t.author_id = p.person_id;

as a more complicated example. Of course, you could wrap all of this in a view for convenience.

**F.45.3.1. Multivalued Results**

The xpath_table function assumes that the results of each XPath query might be multivalued, so
the number of rows returned by the function may not be the same as the number of input documents.
The first row returned contains the first result from each query, the second row the second result from
each query. If one of the queries has fewer values than the others, null values will be returned instead.

In some cases, a user will know that a given XPath query will return only a single result (perhaps
a unique document identifier) — if used alongside an XPath query returning multiple results, the
single-valued result will appear only on the first row of the result. The solution to this is to use the key
field as part of a join against a simpler XPath query. As an example:
CREATE TABLE test (
    id int PRIMARY KEY,
    xml text
);

INSERT INTO test VALUES (1, '<doc num="C1">
    <line num="L1">a1</a><b>2</b><c>3</c></line>
    <line num="L2">a11</a><b>22</b><c>33</c></line>
</doc>');

INSERT INTO test VALUES (2, '<doc num="C2">
    <line num="L1">a111</a><b>222</b><c>333</c></line>
    <line num="L2">a111</a><b>222</b><c>333</c></line>
</doc>');

SELECT * FROM 
    xpath_table('id', 'xml', 'test',
        '/doc/@num|/doc/line/@num|/doc/line/a|/doc/line/b|/
doc/line/c',
        'true')
    AS t(id int, doc_num varchar(10), line_num varchar(10), val1 int, val2 int, val3 int)
WHERE id = 1 ORDER BY doc_num, line_num

id | doc_num | line_num | val1 | val2 | val3 
----+---------+----------+------+------+------
1 | C1      | L1       |    1 |    2 |    3 
1 |         | L2       |   11 |   22 |   33 
(2 rows)

To get doc_num on every line, the solution is to use two invocations of xpath_table and join the results:

SELECT t.*,i.doc_num FROM 
    xpath_table('id', 'xml', 'test',
        '/doc/@num|/doc/line/@num|/doc/line/a|/doc/line/b|/
doc/line/c',
        'true')
    AS t(id int, line_num varchar(10), val1 int, val2 int, val3 int),
    xpath_table('id', 'xml', 'test', '/doc/@num', 'true')
    AS i(id int, doc_num varchar(10))
WHERE i.id=t.id AND i.id=1
ORDER BY doc_num, line_num;

id | line_num | val1 | val2 | val3 | doc_num
----+----------+------+------+------+---------
1 | L1       |    1 |    2 |    3 | C1      
1 | L2       |   11 |   22 |   33 | C1
(2 rows)

F.45.4. XSLT Functions

The following functions are available if libxslt is installed:

F.45.4.1. xslt_process
xslt_process(text document, text stylesheet, text paramlist) returns text

This function applies the XSL stylesheets to the document and returns the transformed result. The paramlist is a list of parameter assignments to be used in the transformation, specified in the form \text{a=1,b=2}. Note that the parameter parsing is very simple-minded: parameter values cannot contain commas!

There is also a two-parameter version of xslt_process which does not pass any parameters to the transformation.

**F.45.5. Author**

John Gray <jgray@azuli.co.uk>

Development of this module was sponsored by Torchbox Ltd. (www.torchbox.com). It has the same BSD license as PostgreSQL.
Appendix G. Additional Supplied Programs

This appendix and the previous one contain information regarding the modules that can be found in the contrib directory of the PostgreSQL distribution. See Appendix F for more information about the contrib section in general and server extensions and plug-ins found in contrib specifically.

This appendix covers utility programs found in contrib. Once installed, either from source or a packaging system, they are found in the bin directory of the PostgreSQL installation and can be used like any other program.

G.1. Client Applications

This section covers PostgreSQL client applications in contrib. They can be run from anywhere, independent of where the database server resides. See also PostgreSQL Client Applications for information about client applications that part of the core PostgreSQL distribution.
**oid2name**

oid2name — resolve OIDs and file nodes in a PostgreSQL data directory

**Synopsis**

`oid2name [option...]`

**Description**

oid2name is a utility program that helps administrators to examine the file structure used by PostgreSQL. To make use of it, you need to be familiar with the database file structure, which is described in Chapter 68.

```
Note

The name “oid2name” is historical, and is actually rather misleading, since most of the time when you use it, you will really be concerned with tables’ filenode numbers (which are the file names visible in the database directories). Be sure you understand the difference between table OIDs and table filenodes!
```

oid2name connects to a target database and extracts OID, filenode, and/or table name information. You can also have it show database OIDs or tablespace OIDs.

**Options**

oid2name accepts the following command-line arguments:

- `-f filenode`
  
  show info for table with filenode *filenode*

- `-i`
  
  include indexes and sequences in the listing

- `-o oid`
  
  show info for table with OID *oid*

- `-q`
  
  omit headers (useful for scripting)

- `-s`
  
  show tablespace OIDs

- `-S`
  
  include system objects (those in information_schema, pg_toast and pg_catalog schemas)

- `-t tablename_pattern`
  
  show info for table(s) matching *tablename_pattern*
Additional Supplied Programs

-\v
--version

Print the oid2name version and exit.

-\x

display more information about each object shown: tablespace name, schema name, and OID

-?  
--help

Show help about oid2name command line arguments, and exit.

oid2name also accepts the following command-line arguments for connection parameters:

-\d database

database to connect to

-\H host

database server's host

-\p port

database server's port

-\U username

user name to connect as

-\P password

password (deprecated — putting this on the command line is a security hazard)

To display specific tables, select which tables to show by using -\o, -\f and/or -\t. -\o takes an OID, -\f takes a filenode, and -\t takes a table name (actually, it's a LIKE pattern, so you can use things like foo%). You can use as many of these options as you like, and the listing will include all objects matched by any of the options. But note that these options can only show objects in the database given by -\d.

If you don't give any of -\o, -\f or -\t, but do give -\d, it will list all tables in the database named by -\d. In this mode, the -\S and -\i options control what gets listed.

If you don’t give -\d either, it will show a listing of database OIDs. Alternatively you can give -\s to get a tablespace listing.

Notes

oid2name requires a running database server with non-corrupt system catalogs. It is therefore of only limited use for recovering from catastrophic database corruption situations.

Examples

$ # what's in this database server, anyway?
$ oid2name
All databases:

<table>
<thead>
<tr>
<th>Oid</th>
<th>Database Name</th>
<th>Tablespace</th>
</tr>
</thead>
<tbody>
<tr>
<td>17228</td>
<td>alvherre</td>
<td>pg_default</td>
</tr>
<tr>
<td>17255</td>
<td>regression</td>
<td>pg_default</td>
</tr>
<tr>
<td>17227</td>
<td>template0</td>
<td>pg_default</td>
</tr>
</tbody>
</table>
Additional Supplied Programs

1      template1  pg_default

$ oid2name -s
All tablespaces:
   Oid  Tablespace Name
----------
  1663   pg_default
  1664   pg_global
 155151  fastdisk
 155152  bigdisk

$ # OK, let's look into database alvherre
$ cd $PGDATA/base/17228

$ # get top 10 db objects in the default tablespace, ordered by size
$ ls -lS * | head -10
-rw-------  1 alvherre alvherre 136536064 sep 14 09:51 155173
-rw-------  1 alvherre alvherre 17965056 sep 14 09:51 1155291
-rw-------  1 alvherre alvherre 1204224 sep 14 09:51 16717
-rw-------  1 alvherre alvherre  581632 sep  6 17:51 1255
-rw-------  1 alvherre alvherre  237568 sep 14 09:50 16674
-rw-------  1 alvherre alvherre  212992 sep 14 09:51 1249
-rw-------  1 alvherre alvherre  204800 sep 14 09:51 16684
-rw-------  1 alvherre alvherre  196608 sep 14 09:50 16700
-rw-------  1 alvherre alvherre  163840 sep 14 09:50 16699
-rw-------  1 alvherre alvherre  122880 sep  6 17:51 16751

$ # I wonder what file 155173 is ...
$ oid2name -d alvherre -f 155173
From database "alvherre":
   Filenode Table Name
----------------------
  155173       accounts

$ # you can ask for more than one object
$ oid2name -d alvherre -f 155173 -f 1155291
From database "alvherre":
   Filenode Table Name
-------------------------
  155173       accounts
 1155291  accounts_pkey

$ # you can mix the options, and get more details with -x
$ oid2name -d alvherre -t accounts -f 1155291 -x
From database "alvherre":
   Filenode Table Name    Oid   Schema  Tablespace
--------------------------------------------------------
  155173       accounts  155173  public  pg_default
 1155291  accounts_pkey 1155291  public  pg_default

$ # show disk space for every db object
$ du [0-9]* |
> while read SIZE FILENODE
> do
>   echo "$SIZE  'oid2name -q -d alvherre -i -f $FILENODE'"
> done
16   1155287  branches_pkey
Additional Supplied Programs

16 1155289 tellers_pkey
17561 1155291 accounts_pkey
...

$ # same, but sort by size
$ du [0-9]* | sort -rn | while read SIZE FN
> do
>   echo "$SIZE `oid2name -q -d alvherre -f $FN`"
> done
133466 155173 accounts
17561 1155291 accounts_pkey
1177 16717 pg_proc_proname_args_nsp_index
...

$ # If you want to see what's in tablespaces, use the pg_tblspc directory
$ cd $PGDATA/pg_tblspc
$ oid2name -s
All tablespaces:
  Oid  Tablespace Name
  -------------------------
    1663     pg_default
    1664     pg_global
  155151    fastdisk
  155152    bigdisk

$ # what databases have objects in tablespace "fastdisk"?
$ ls -d 155151/*
155151/17228/  155151/PG_VERSION

$ # Oh, what was database 17228 again?
$ oid2name
All databases:
  Oid  Database Name  Tablespace
  ----------------------------------
  17228       alvherre  pg_default
  17255     regression  pg_default
  17227      template0  pg_default
       1      template1  pg_default

$ # Let's see what objects does this database have in the tablespace.
$ cd 155151/17228
$ ls -l
total 0
-rw------- 1 postgres postgres 0 sep 13 23:20 155156

$ # OK, this is a pretty small table ... but which one is it?
$ oid2name -d alvherre -f 155156
From database "alvherre":
  Filenode  Table Name
  ----------------------
    155156    foo

Author

B. Palmer <bpalmer@crimelabs.net>
vacuumlo

vacuumlo — remove orphaned large objects from a PostgreSQL database

Synopsis

vacuumlo [option...] dbname...

Description

vacuumlo is a simple utility program that will remove any “orphaned” large objects from a PostgreSQL database. An orphaned large object (LO) is considered to be any LO whose OID does not appear in any oid or lo data column of the database.

If you use this, you may also be interested in the lo_manage trigger in the lo module. lo_manage is useful to try to avoid creating orphaned LOs in the first place.

All databases named on the command line are processed.

Options

vacuumlo accepts the following command-line arguments:

- \( l \) limit

Remove no more than \( limit \) large objects per transaction (default 1000). Since the server acquires a lock per LO removed, removing too many LOs in one transaction risks exceeding max_locks_per_transaction. Set the limit to zero if you want all removals done in a single transaction.

- \( n \)

Don’t remove anything, just show what would be done.

- \( v \)

Write a lot of progress messages.

- \( V \)

--version

Print the vacuumlo version and exit.

- \( ? \)

--help

Show help about vacuumlo command line arguments, and exit.

vacuumlo also accepts the following command-line arguments for connection parameters:

- \( h \) hostname

Database server’s host.

- \( p \) port

Database server’s port.
-U username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W

Force vacuumlo to prompt for a password before connecting to a database.

This option is never essential, since vacuumlo will automatically prompt for a password if the server demands password authentication. However, vacuumlo will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

Notes

vacuumlo works by the following method: First, vacuumlo builds a temporary table which contains all of the OIDs of the large objects in the selected database. It then scans through all columns in the database that are of type oid or lo, and removes matching entries from the temporary table. (Note: Only types with these names are considered; in particular, domains over them are not considered.) The remaining entries in the temporary table identify orphaned LOs. These are removed.

Author

Peter Mount <peter@retep.org.uk>

G.2. Server Applications

This section covers PostgreSQL server-related applications in contrib. They are typically run on the host where the database server resides. See also PostgreSQL Server Applications for information about server applications that part of the core PostgreSQL distribution.
Additional Supplied Programs

pg_standby

pg_standby — supports the creation of a PostgreSQL warm standby server

Synopsis

pg_standby [option...] archivelocation nextwalfile walfilepath [restart-walfile]

Description

pg_standby supports creation of a “warm standby” database server. It is designed to be a production-ready program, as well as a customizable template should you require specific modifications.

pg_standby is designed to be a waiting restore_command, which is needed to turn a standard archive recovery into a warm standby operation. Other configuration is required as well, all of which is described in the main server manual (see Section 26.2).

To configure a standby server to use pg_standby, put this into its recovery.conf configuration file:

restore_command = 'pg_standby archiveDir %f %p %r'

where archiveDir is the directory from which WAL segment files should be restored.

If restartwalfile is specified, normally by using the %r macro, then all WAL files logically preceding this file will be removed from archivelocation. This minimizes the number of files that need to be retained, while preserving crash-restart capability. Use of this parameter is appropriate if the archivelocation is a transient staging area for this particular standby server, but not when the archivelocation is intended as a long-term WAL archive area.

pg_standby assumes that archivelocation is a directory readable by the server-owning user. If restartwalfile (or -k) is specified, the archivelocation directory must be writable too.

There are two ways to fail over to a “warm standby” database server when the master server fails:

Smart Failover

In smart failover, the server is brought up after applying all WAL files available in the archive. This results in zero data loss, even if the standby server has fallen behind, but if there is a lot of unapplied WAL it can be a long time before the standby server becomes ready. To trigger a smart failover, create a trigger file containing the word smart, or just create it and leave it empty.

Fast Failover

In fast failover, the server is brought up immediately. Any WAL files in the archive that have not yet been applied will be ignored, and all transactions in those files are lost. To trigger a fast failover, create a trigger file and write the word fast into it. pg_standby can also be configured to execute a fast failover automatically if no new WAL file appears within a defined interval.

Options

pg_standby accepts the following command-line arguments:

- Use cp or copy command to restore WAL files from archive. This is the only supported behavior so this option is useless.
-d

Print lots of debug logging output on stderr.

-k

Remove files from archive location so that no more than this many WAL files before the current one are kept in the archive. Zero (the default) means not to remove any files from archive location. This parameter will be silently ignored if restartwalsfile is specified, since that specification method is more accurate in determining the correct archive cut-off point. Use of this parameter is deprecated as of PostgreSQL 8.3; it is safer and more efficient to specify a restartwalsfile parameter. A too small setting could result in removal of files that are still needed for a restart of the standby server, while a too large setting wastes archive space.

-r maxretries

Set the maximum number of times to retry the copy command if it fails (default 3). After each failure, we wait for sleeptime * num_retries so that the wait time increases progressively. So by default, we will wait 5 secs, 10 secs, then 15 secs before reporting the failure back to the standby server. This will be interpreted as end of recovery and the standby will come up fully as a result.

-s sleeptime

Set the number of seconds (up to 60, default 5) to sleep between tests to see if the WAL file to be restored is available in the archive yet. The default setting is not necessarily recommended; consult Section 26.2 for discussion.

-t triggerfile

Specify a trigger file whose presence should cause failover. It is recommended that you use a structured file name to avoid confusion as to which server is being triggered when multiple servers exist on the same system; for example /tmp/pgsql.trigger.5432.

-V

--version

Print the pg_standby version and exit.

-w maxwaittime

Set the maximum number of seconds to wait for the next WAL file, after which a fast failover will be performed. A setting of zero (the default) means wait forever. The default setting is not necessarily recommended; consult Section 26.2 for discussion.

-?

--help

Show help about pg_standby command line arguments, and exit.

Notes

pg_standby is designed to work with PostgreSQL 8.2 and later.

PostgreSQL 8.3 provides the %r macro, which is designed to let pg_standby know the last file it needs to keep. With PostgreSQL 8.2, the --k option must be used if archive cleanup is required. This option remains available in 8.3, but its use is deprecated.

PostgreSQL 8.4 provides the recovery_end_command option. Without this option a leftover trigger file can be hazardous.
pg_standby is written in C and has an easy-to-modify source code, with specifically designated sections to modify for your own needs.

Examples

On Linux or Unix systems, you might use:

archive_command = 'cp %p .../archive/%f'

restore_command = 'pg_standby -d -s 2 -t /tmp/pgsql.trigger.5442 .../archive %f %p %r 2>>standby.log'

recovery_end_command = 'rm -f /tmp/pgsql.trigger.5442'

where the archive directory is physically located on the standby server, so that the archive_command is accessing it across NFS, but the files are local to the standby (enabling use of ln). This will:

• produce debugging output in standby.log
• sleep for 2 seconds between checks for next WAL file availability
• stop waiting only when a trigger file called /tmp/pgsql.trigger.5442 appears, and perform failover according to its content
• remove the trigger file when recovery ends
• remove no-longer-needed files from the archive directory

On Windows, you might use:

archive_command = 'copy %p ...\archive\%f'

restore_command = 'pg_standby -d -s 5 -t C:\pgsql.trigger.5442 ...\archive %f %p %r 2>>standby.log'

recovery_end_command = 'del C:\pgsql.trigger.5442'

Note that backslashes need to be doubled in the archive_command, but not in the restore_command or recovery_end_command. This will:

• use the copy command to restore WAL files from archive
• produce debugging output in standby.log
• sleep for 5 seconds between checks for next WAL file availability
• stop waiting only when a trigger file called C:\pgsql.trigger.5442 appears, and perform failover according to its content
• remove the trigger file when recovery ends
• remove no-longer-needed files from the archive directory

The copy command on Windows sets the final file size before the file is completely copied, which would ordinarily confuse pg_standby. Therefore pg_standby waits sleeptime seconds once it sees the proper file size. GNUWin32's cp sets the file size only after the file copy is complete.

Since the Windows example uses copy at both ends, either or both servers might be accessing the archive directory across the network.
Author

Simon Riggs <simon@2ndquadrant.com>

See Also

pg_archivecleanup
Appendix H. External Projects

PostgreSQL is a complex software project, and managing the project is difficult. We have found that many enhancements to PostgreSQL can be more efficiently developed separately from the core project.

H.1. Client Interfaces

There are only two client interfaces included in the base PostgreSQL distribution:

- `libpq` is included because it is the primary C language interface, and because many other client interfaces are built on top of it.

- ECPG is included because it depends on the server-side SQL grammar, and is therefore sensitive to changes in PostgreSQL itself.

All other language interfaces are external projects and are distributed separately. Table H.1 includes a list of some of these projects. Note that some of these packages might not be released under the same license as PostgreSQL. For more information on each language interface, including licensing terms, refer to its website and documentation.

Table H.1. Externally Maintained Client Interfaces

<table>
<thead>
<tr>
<th>Name</th>
<th>Language</th>
<th>Comments</th>
<th>Website</th>
</tr>
</thead>
<tbody>
<tr>
<td>DBD::Pg</td>
<td>Perl</td>
<td>Perl DBI driver</td>
<td><a href="https://metacpan.org/release/DBD-Pg">https://metacpan.org/release/DBD-Pg</a></td>
</tr>
<tr>
<td>JDBC</td>
<td>Java</td>
<td>Type 4 JDBC driver</td>
<td><a href="https://jdbc.postgresql.org/">https://jdbc.postgresql.org/</a></td>
</tr>
<tr>
<td>libpqxx</td>
<td>C++</td>
<td>C++ interface</td>
<td><a href="http://pqxx.org/">http://pqxx.org/</a></td>
</tr>
<tr>
<td>node-postgres</td>
<td>JavaScript</td>
<td>Node.js driver</td>
<td><a href="https://node-postgres.com/">https://node-postgres.com/</a></td>
</tr>
<tr>
<td>Npgsql</td>
<td>.NET</td>
<td>.NET data provider</td>
<td><a href="http://www.npgsql.org/">http://www.npgsql.org/</a></td>
</tr>
<tr>
<td>pgtcl</td>
<td>Tcl</td>
<td></td>
<td><a href="https://github.com/flightaware/Pgtcl">https://github.com/flightaware/Pgtcl</a></td>
</tr>
<tr>
<td>pgtclng</td>
<td>Tcl</td>
<td></td>
<td><a href="https://sourceforge.net/projects/pgtclng/">https://sourceforge.net/projects/pgtclng/</a></td>
</tr>
<tr>
<td>pq</td>
<td>Go</td>
<td>Pure Go driver for Go's database/sql</td>
<td><a href="https://github.com/libpq">https://github.com/libpq</a></td>
</tr>
<tr>
<td>psqlODBC</td>
<td>ODBC</td>
<td>ODBC driver</td>
<td><a href="https://odbc.postgresql.org/">https://odbc.postgresql.org/</a></td>
</tr>
<tr>
<td>psycopg</td>
<td>Python</td>
<td>DB API 2.0-compliant</td>
<td><a href="http://initd.org/psycopg/">http://initd.org/psycopg/</a></td>
</tr>
</tbody>
</table>

H.2. Administration Tools

There are several administration tools available for PostgreSQL. The most popular is pgAdmin¹, and there are several commercially available ones as well.

¹ https://www.pgadmin.org/
H.3. Procedural Languages

PostgreSQL includes several procedural languages with the base distribution: PL/pgSQL, PL/Tcl, PL/Perl, and PL/Python.

In addition, there are a number of procedural languages that are developed and maintained outside the core PostgreSQL distribution. Table H.2 lists some of these packages. Note that some of these projects might not be released under the same license as PostgreSQL. For more information on each procedural language, including licensing information, refer to its website and documentation.

Table H.2. Externally Maintained Procedural Languages

<table>
<thead>
<tr>
<th>Name</th>
<th>Language</th>
<th>Website</th>
</tr>
</thead>
<tbody>
<tr>
<td>PL/Java</td>
<td>Java</td>
<td><a href="https://tada.github.io/pljava/">https://tada.github.io/pljava/</a></td>
</tr>
<tr>
<td>PL/Lua</td>
<td>Lua</td>
<td><a href="https://github.com/pllua/pllua">https://github.com/pllua/pllua</a></td>
</tr>
<tr>
<td>PL/R</td>
<td>R</td>
<td><a href="http://www.joeconway.com/plr.html">http://www.joeconway.com/plr.html</a></td>
</tr>
<tr>
<td>PL/sh</td>
<td>Unix shell</td>
<td><a href="https://github.com/petere/plsh">https://github.com/petere/plsh</a></td>
</tr>
<tr>
<td>PL/v8</td>
<td>JavaScript</td>
<td><a href="https://github.com/plv8/plv8">https://github.com/plv8/plv8</a></td>
</tr>
</tbody>
</table>

H.4. Extensions

PostgreSQL is designed to be easily extensible. For this reason, extensions loaded into the database can function just like features that are built in. The contrib/ directory shipped with the source code contains several extensions, which are described in Appendix F. Other extensions are developed independently, like PostGIS\(^2\). Even PostgreSQL replication solutions can be developed externally. For example, Slony-I\(^3\) is a popular master/standby replication solution that is developed independently from the core project.

---

\(^2\) http://postgis.net/
\(^3\) http://www.slony.info
Appendix I. The Source Code Repository

The PostgreSQL source code is stored and managed using the Git version control system. A public mirror of the master repository is available; it is updated within a minute of any change to the master repository.

Our wiki, https://wiki.postgresql.org/wiki/Working_with_Git, has some discussion on working with Git.

Note that building PostgreSQL from the source repository requires reasonably up-to-date versions of bison, flex, and Perl. These tools are not needed to build from a distribution tarball, because the files that these tools are used to build are included in the tarball. Other tool requirements are the same as shown in Section 16.2.

I.1. Getting The Source via Git

With Git you will make a copy of the entire code repository on your local machine, so you will have access to all history and branches offline. This is the fastest and most flexible way to develop or test patches.

Git

1. You will need an installed version of Git, which you can get from https://git-scm.com. Many systems already have a recent version of Git installed by default, or available in their package distribution system.

2. To begin using the Git repository, make a clone of the official mirror:

```bash
git clone https://git.postgresql.org/git/postgresql.git
```

This will copy the full repository to your local machine, so it may take a while to complete, especially if you have a slow Internet connection. The files will be placed in a new subdirectory `postgresql` of your current directory.

The Git mirror can also be reached via the Git protocol. Just change the URL prefix to `git`, as in:

```bash
git clone git://git.postgresql.org/git/postgresql.git
```

3. Whenever you want to get the latest updates in the system, `cd` into the repository, and run:

```bash
git fetch
```

Git can do a lot more things than just fetch the source. For more information, consult the Git man pages, or see the website at https://git-scm.com.
Appendix J. Documentation

PostgreSQL has four primary documentation formats:

• Plain text, for pre-installation information
• HTML, for on-line browsing and reference
• PDF, for printing
• man pages, for quick reference.

Additionally, a number of plain-text README files can be found throughout the PostgreSQL source tree, documenting various implementation issues.

HTML documentation and man pages are part of a standard distribution and are installed by default. PDF format documentation is available separately for download.

J.1. DocBook

The documentation sources are written in DocBook, which is a markup language defined in XML. In what follows, the terms DocBook and XML are both used, but technically they are not interchangeable.

DocBook allows an author to specify the structure and content of a technical document without worrying about presentation details. A document style defines how that content is rendered into one of several final forms. DocBook is maintained by the OASIS group. The official DocBook site has good introductory and reference documentation and a complete O'Reilly book for your online reading pleasure. The NewbieDoc Docbook Guide is very helpful for beginners. The FreeBSD Documentation Project also uses DocBook and has some good information, including a number of style guidelines that might be worth considering.

J.2. Tool Sets

The following tools are used to process the documentation. Some might be optional, as noted.

DocBook DTD

This is the definition of DocBook itself. We currently use version 4.2; you cannot use later or earlier versions. You need the XML variant of the DocBook DTD, not the SGML variant.

DocBook XSL Stylesheets

These contain the processing instructions for converting the DocBook sources to other formats, such as HTML.

The minimum required version is currently 1.77.0, but it is recommended to use the latest available version for best results.

Libxml2 for xmllint

This library and the xmllint tool it contains are used for processing XML. Many developers will already have Libxml2 installed, because it is also used when building the PostgreSQL code. Note, however, that xmllint might need to be installed from a separate subpackage.

---

1 https://www.oasis-open.org
2 https://www.oasis-open.org/docbook/
4 https://www.freebsd.org/docproj/docproj.html
5 http://www.oasis-open.org/docbook/
7 http://xmlsoft.org/
Libxslt\(^8\) for \texttt{xsltproc}

\texttt{xsltproc} is an XSLT processor, that is, a program to convert XML to other formats using XSLT stylesheets.

FOP\(^9\)

This is a program for converting, among other things, XML to PDF.

We have documented experience with several installation methods for the various tools that are needed to process the documentation. These will be described below. There might be some other packaged distributions for these tools. Please report package status to the documentation mailing list, and we will include that information here.

You can get away with not installing DocBook XML and the DocBook XSLT stylesheets locally, because the required files will be downloaded from the Internet and cached locally. This may in fact be the preferred solution if your operating system packages provide only an old version of especially the stylesheets or if no packages are available at all. See the \texttt{--nonet} option for \texttt{xmllint} and \texttt{xsltproc} for more information.

\section*{J.2.1. Installation on Fedora, RHEL, and Derivatives}

To install the required packages, use:

\begin{verbatim}
yum install docbook-dtds docbook-style-xsl fop libxslt
\end{verbatim}

\section*{J.2.2. Installation on FreeBSD}

To install the required packages with \texttt{pkg}, use:

\begin{verbatim}
pkg install docbook-xml docbook-xsl fop libxslt
\end{verbatim}

When building the documentation from the \texttt{doc} directory you'll need to use \texttt{gmake}, because the makefile provided is not suitable for FreeBSD's \texttt{make}.

\section*{J.2.3. Debian Packages}

There is a full set of packages of the documentation tools available for Debian GNU/Linux. To install, simply use:

\begin{verbatim}
apt-get install docbook-xml docbook-xsl fop libxml2-utils xsltproc
\end{verbatim}

\section*{J.2.4. macOS}

On macOS, you can build the HTML and man documentation without installing anything extra. If you want to build PDFs or want to install a local copy of DocBook, you can get those from your preferred package manager.

If you use MacPorts, the following will get you set up:

\begin{verbatim}
sudo port install docbook-xml-4.2 docbook-xsl fop
\end{verbatim}

If you use Homebrew, use this:

\begin{verbatim}
sudo brew install docbook-xml docbook-xsl fop
\end{verbatim}

\footnote{8 \url{http://xmlsoft.org/XSLT/}}
\footnote{9 \url{https://xmlgraphics.apache.org/fop/}}

3471
brew install docbook docbook-xsl fop

J.2.5. Detection by configure

Before you can build the documentation you need to run the configure script as you would when building the PostgreSQL programs themselves. Check the output near the end of the run, it should look something like this:

```
checking for xmllint... xmllint
checking for DocBook XML V4.2... yes
checking for dbtoepub... dbtoepub
checking for xsltproc... xsltproc
checking for fop... fop
```

If xmllint was not found then some of the following tests will be skipped.

J.3. Building The Documentation

Once you have everything set up, change to the directory `doc/src/sgml` and run one of the commands described in the following subsections to build the documentation. (Remember to use GNU make.)

J.3.1. HTML

To build the HTML version of the documentation:

```
doc/src/sgml$ make html
```

This is also the default target. The output appears in the subdirectory `html`.

To produce HTML documentation with the stylesheet used on postgresql.org instead of the default simple style use:

```
doc/src/sgml$ make STYLE=website html
```

J.3.2. Manpages

We use the DocBook XSL stylesheets to convert DocBook refentry pages to *roff output suitable for man pages. To create the man pages, use the command:

```
doc/src/sgml$ make man
```

J.3.3. PDF

To produce a PDF rendition of the documentation using FOP, you can use one of the following commands, depending on the preferred paper format:

- For A4 format:

```
doc/src/sgml$ make postgres-A4.pdf
```

10 https://www.postgresql.org/docs/current/
• For U.S. letter format:

```
make postgres-US.pdf
```

Because the PostgreSQL documentation is fairly big, FOP will require a significant amount of memory. Because of that, on some systems, the build will fail with a memory-related error message. This can usually be fixed by configuring Java heap settings in the configuration file `~/.foprc`, for example:

```
# FOP binary distribution
FOP_OPTS='-Xmx1500m'
# Debian
JAVA_ARGS='-Xmx1500m'
# Red Hat
ADDITIONAL_FLAGS='-Xmx1500m'
```

There is a minimum amount of memory that is required, and to some extent more memory appears to make things a bit faster. On systems with very little memory (less than 1 GB), the build will either be very slow due to swapping or will not work at all.

Other XSL-FO processors can also be used manually, but the automated build process only supports FOP.

### J.3.4. Plain Text Files

The installation instructions are also distributed as plain text, in case they are needed in a situation where better reading tools are not available. The `INSTALL` file corresponds to Chapter 16, with some minor changes to account for the different context. To recreate the file, change to the directory `doc/src/sgml` and enter `make INSTALL`.

In the past, the release notes and regression testing instructions were also distributed as plain text, but this practice has been discontinued.

### J.3.5. Syntax Check

Building the documentation can take very long. But there is a method to just check the correct syntax of the documentation files, which only takes a few seconds:

```
make check
```

### J.4. Documentation Authoring

The documentation sources are most conveniently modified with an editor that has a mode for editing XML, and even more so if it has some awareness of XML schema languages so that it can know about DocBook syntax specifically.

Note that for historical reasons the documentation source files are named with an extension `.sgml` even though they are now XML files. So you might need to adjust your editor configuration to set the correct mode.

### J.4.1. Emacs

nXML Mode, which ships with Emacs, is the most common mode for editing XML documents with Emacs. It will allow you to use Emacs to insert tags and check markup consistency, and it supports DocBook out of the box. Check the nXML manual\(^\text{11}\) for detailed documentation.

src/tools/editors/emacs.samples contains recommended settings for this mode.

J.5. Style Guide

J.5.1. Reference Pages

Reference pages should follow a standard layout. This allows users to find the desired information more quickly, and it also encourages writers to document all relevant aspects of a command. Consistency is not only desired among PostgreSQL reference pages, but also with reference pages provided by the operating system and other packages. Hence the following guidelines have been developed. They are for the most part consistent with similar guidelines established by various operating systems.

Reference pages that describe executable commands should contain the following sections, in this order. Sections that do not apply can be omitted. Additional top-level sections should only be used in special circumstances; often that information belongs in the “Usage” section.

Name

This section is generated automatically. It contains the command name and a half-sentence summary of its functionality.

Synopsis

This section contains the syntax diagram of the command. The synopsis should normally not list each command-line option; that is done below. Instead, list the major components of the command line, such as where input and output files go.

Description

Several paragraphs explaining what the command does.

Options

A list describing each command-line option. If there are a lot of options, subsections can be used.

Exit Status

If the program uses 0 for success and non-zero for failure, then you do not need to document it. If there is a meaning behind the different non-zero exit codes, list them here.

Usage

Describe any sublanguage or run-time interface of the program. If the program is not interactive, this section can usually be omitted. Otherwise, this section is a catch-all for describing run-time features. Use subsections if appropriate.

Environment

List all environment variables that the program might use. Try to be complete; even seemingly trivial variables like SHELL might be of interest to the user.

Files

List any files that the program might access implicitly. That is, do not list input and output files that were specified on the command line, but list configuration files, etc.

Diagnostics

Explain any unusual output that the program might create. Refrain from listing every possible error message. This is a lot of work and has little use in practice. But if, say, the error messages have a standard format that the user can parse, this would be the place to explain it.
Documentation

Notes

Anything that doesn't fit elsewhere, but in particular bugs, implementation flaws, security considerations, compatibility issues.

Examples

Examples

History

If there were some major milestones in the history of the program, they might be listed here. Usually, this section can be omitted.

Author

Author (only used in the contrib section)

See Also

Cross-references, listed in the following order: other PostgreSQL command reference pages, PostgreSQL SQL command reference pages, citation of PostgreSQL manuals, other reference pages (e.g., operating system, other packages), other documentation. Items in the same group are listed alphabetically.

Reference pages describing SQL commands should contain the following sections: Name, Synopsis, Description, Parameters, Outputs, Notes, Examples, Compatibility, History, See Also. The Parameters section is like the Options section, but there is more freedom about which clauses of the command can be listed. The Outputs section is only needed if the command returns something other than a default command-completion tag. The Compatibility section should explain to what extent this command conforms to the SQL standard(s), or to which other database system it is compatible. The See Also section of SQL commands should list SQL commands before cross-references to programs.
Appendix K. Acronyms

This is a list of acronyms commonly used in the PostgreSQL documentation and in discussions about PostgreSQL.

ANSI
    American National Standards Institute¹

API
    Application Programming Interface²

ASCII
    American Standard Code for Information Interchange³

BKI
    Backend Interface

CA
    Certificate Authority⁴

CIDR
    Classless Inter-Domain Routing⁵

CPAN
    Comprehensive Perl Archive Network⁶

CRL
    Certificate Revocation List⁷

CSV
    Comma Separated Values⁸

CTE
    Common Table Expression

CVE
    Common Vulnerabilities and Exposures⁹

DBA
    Database Administrator¹⁰

¹ https://en.wikipedia.org/wiki/American_National_Standards_Institute
³ https://en.wikipedia.org/wiki/Ascii
⁵ https://en.wikipedia.org/wiki/Classless_Inter-Domain_Routing
⁶ https://www.cpan.org/
⁸ https://en.wikipedia.org/wiki/Comma-separated_values
⁹ http://cve.mitre.org/
¹⁰ https://en.wikipedia.org/wiki/Database_administrator
Acronyms

DBI
Database Interface (Perl)

DBMS
Database Management System

DDL
Data Definition Language, SQL commands such as CREATE TABLE, ALTER USER

DML
Data Manipulation Language, SQL commands such as INSERT, UPDATE, DELETE

DST
Daylight Saving Time

ECPG
Embedded C for PostgreSQL

ESQL
Embedded SQL

FAQ
Frequently Asked Questions

FSM
Free Space Map

GEQO
Genetic Query Optimizer

GIN
Generalized Inverted Index

GiST
Generalized Search Tree

Git

GMT
Greenwich Mean Time

11 https://dbi.perl.org/
12 https://en.wikipedia.org/wiki/Dbms
16 https://en.wikipedia.org/wiki/Embedded_SQL
17 https://en.wikipedia.org/wiki/FAQ
18 https://en.wikipedia.org/wiki/Git_(software)
Acronyms

GSSAPI

Generic Security Services Application Programming Interface\(^{20}\)

GUC

Grand Unified Configuration, the PostgreSQL subsystem that handles server configuration

HBA

Host-Based Authentication

HOT

Heap-Only Tuples\(^{21}\)

IEC

International Electrotechnical Commission\(^{22}\)

IEEE

Institute of Electrical and Electronics Engineers\(^{23}\)

IPC

Inter-Process Communication\(^{24}\)

ISO

International Organization for Standardization\(^{25}\)

ISSN

International Standard Serial Number\(^{26}\)

JDBC

Java Database Connectivity\(^{27}\)

JIT

Just-in-Time compilation\(^{28}\)

JSON

JavaScript Object Notation\(^{29}\)

LDAP

Lightweight Directory Access Protocol\(^{30}\)

\(^{20}\) https://en.wikipedia.org/wiki/Generic_Security_Services_Application_Programming_Interface
\(^{21}\) https://git.postgresql.org/gitweb/?p=postgresql.git;a=blt;f=src/backend/access/heap/README.HOT;hb=HEAD
\(^{22}\) https://en.wikipedia.org/wiki/International_Electrotechnical_Commission
\(^{23}\) http://standards.ieee.org/
\(^{24}\) https://en.wikipedia.org/wiki/Inter-process_communication
\(^{25}\) https://www.iso.org/home.html
\(^{26}\) https://en.wikipedia.org/wiki/Issn
\(^{27}\) https://en.wikipedia.org/wiki/Java_Database_Connectivity
\(^{29}\) http://json.org
Acronyms

LSN
Log Sequence Number, see pg_lsn and WAL Internals.

MSVC
Microsoft Visual C

MVCC
Multi-Version Concurrency Control

NLS
National Language Support

ODBC
Open Database Connectivity

OID
Object Identifier

OLAP
Online Analytical Processing

OLTP
Online Transaction Processing

ORDBMS
Object-Relational Database Management System

PAM
Pluggable Authentication Modules

PGSQL
PostgreSQL

PGXS
PostgreSQL Extension System

PID
Process Identifier

PITR
Point-In-Time Recovery (Continuous Archiving)

33 https://en.wikipedia.org/wiki/Open_Database_Connectivity
34 https://en.wikipedia.org/wiki/Olap
35 https://en.wikipedia.org/wiki/OLTP
36 https://en.wikipedia.org/wiki/ORDBMS
37 https://en.wikipedia.org/wiki/Pluggable_Authentication_Modules
38 https://en.wikipedia.org/wiki/Process_identifier
Acronyms

PL
Procedural Languages (server-side)

POSIX
Portable Operating System Interface\(^{39}\)

RDBMS
Relational Database Management System\(^{40}\)

RFC
Request For Comments\(^{41}\)

SGML
Standard Generalized Markup Language\(^{42}\)

SPI
Server Programming Interface

SP-GiST
Space-Partitioned Generalized Search Tree

SQL
Structured Query Language\(^{43}\)

SRF
Set-Returning Function

SSH
Secure Shell\(^{44}\)

SSL
Secure Sockets Layer\(^{45}\)

SSPI
Security Support Provider Interface\(^{46}\)

SYSV
Unix System V\(^{47}\)

TCP/IP
Transmission Control Protocol (TCP) / Internet Protocol (IP)\(^{48}\)

\(^{39}\) https://en.wikipedia.org/wiki/POSIX

\(^{40}\) https://en.wikipedia.org/wiki/Relational_database_management_system

\(^{41}\) https://en.wikipedia.org/wiki/Request_for_Comments

\(^{42}\) https://en.wikipedia.org/wiki/SGML

\(^{43}\) https://en.wikipedia.org/wiki/SQL

\(^{44}\) https://en.wikipedia.org/wiki/Secure_Shell

\(^{45}\) https://en.wikipedia.org/wiki/Secure_Sockets_Layer


\(^{47}\) https://en.wikipedia.org/wiki/System_V

Acronyms

TID
TUPLE IDENTIFIER

TOAST
THE OVERSIZED-ATTRIBUTE STORAGE TECHNIQUE

TPC
TRANSACTION PROCESSING PERFORMANCE COUNCIL

URL
UNIFORM RESOURCE LOCATOR

UTC
COORDINATED UNIVERSAL TIME

UTF
UNICODE TRANSFORMATION FORMAT

UTF8
EIGHT-BIT UNICODE TRANSFORMATION FORMAT

UUID
UNIVERSALLY UNIQUE IDENTIFIER

WAL
WRITE-AHEAD LOG

XID
TRANSACTION IDENTIFIER

XML
EXTENSIBLE MARKUP LANGUAGE

49 http://www.tpc.org/
50 https://en.wikipedia.org/wiki/URL
52 http://www.unicode.org/
53 https://en.wikipedia.org/wiki/Utf8
Bibliography

Selected references and readings for SQL and PostgreSQL.

Some white papers and technical reports from the original POSTGRES development team are available at the University of California, Berkeley, Computer Science Department web site¹.

**SQL Reference Books**


**PostgreSQL-specific Documentation**


[fong] *The design and implementation of the POSTGRES query optimizer*². Zelaine Fong. University of California, Berkeley, Computer Science Department.

**Proceedings and Articles**


¹ http://db.cs.berkeley.edu/papers/
² http://db.cs.berkeley.edu/papers/UCB-MS-zfong.pdf
⁴ http://citeseer.ist.psu.edu/seshadri95generalized.html


6 http://db.cs.berkeley.edu/papers/ERL-M87-06.pdf
7 http://db.cs.berkeley.edu/papers/ERL-M89-82.pdf
9 http://db.cs.berkeley.edu/papers/ERL-M90-34.pdf
Index

Symbols
S, 42
$libdir, 1050
$libdir/plugins, 582, 1728
*, 120
.pgpass, 831
.pg_service.conf, 831
::, 49
_PG_fini, 1049
_PG_init, 1049
_PG_output_plugin_init, 1336

A
abbrev, 263
ABORT, 1352
abs, 202
acos, 204
acosd, 204
administration tools
externally maintained, 3467
adminpack, 3294
advisory lock, 436
age, 245
aggregate function, 13
built-in, 303
invocation, 44
moving aggregate, 1072
ordered set, 1075
partial aggregation, 1076
polymorphic, 1073
support functions for, 1077
user-defined, 1070
variadic, 1073
AIX
installation on, 490
IPC configuration, 511
akeys, 3355
alias
for table name in query, 13
in the FROM clause, 110
in the select list, 121
ALL, 312, 315
allow_system_table_mods configuration parameter, 589
ALTER AGGREGATE, 1353
ALTER COLLATION, 1355
ALTER CONVERSION, 1357
ALTER DATABASE, 1359
ALTER DEFAULT PRIVILEGES, 1362
ALTER DOMAIN, 1365
ALTER EVENT TRIGGER, 1368
ALTER EXTENSION, 1369
ALTER FOREIGN DATA WRAPPER, 1373
ALTER FOREIGN TABLE, 1375
ALTER FUNCTION, 1380
ALTER GROUP, 1384
ALTER INDEX, 1386
ALTER LANGUAGE, 1389
ALTER LARGE OBJECT, 1390
ALTER MATERIALIZED VIEW, 1391
ALTER OPERATOR, 1393
ALTER OPERATOR CLASS, 1395
ALTER OPERATOR FAMILY, 1396
ALTER POLICY, 1400
ALTER PROCEDURE, 1402
ALTER PUBLICATION, 1405
ALTER ROLE, 616, 1407
ALTER ROUTINE, 1411
ALTER RULE, 1413
ALTER SCHEMA, 1414
ALTER SEQUENCE, 1415
ALTER SERVER, 1418
ALTER STATISTICS, 1420
ALTER SUBSCRIPTION, 1421
ALTER SYSTEM, 1423
ALTER TABLE, 1425
ALTER TABLESPACE, 1441
ALTER TEXT SEARCH CONFIGURATION, 1443
ALTER TEXT SEARCH DICTIONARY, 1445
ALTER TEXT SEARCH PARSER, 1447
ALTER TEXT SEARCH TEMPLATE, 1448
ALTER TRIGGER, 1449
ALTER TYPE, 1451
ALTER USER, 1455
ALTER USER MAPPING, 1456
ALTER VIEW, 1457
amcheck, 3295
ANALYZE, 644, 1459
AND (operator), 198
anonymous code blocks, 1646
any, 196
ANY, 305, 312, 315
anyarray, 196
anyelement, 196
anyenum, 196
anyonarray, 196
anyrange, 196
applicable role, 970
application_name configuration parameter, 567
arbitrary precision numbers, 133
archive_cleanup_command recovery parameter, 690
archive_command configuration parameter, 550
archive_mode configuration parameter, 550
archive_timeout configuration parameter, 551
area, 259
armor, 3388
array, 172
accessing, 174
constant, 173
constructor, 51
declaration, 172
I/O, 180

3484
modifying, 176
of user-defined type, 1080
searching, 179
ARRAY, 51
determination of result type, 368
array_agg, 304, 3359
array_append, 299
array_cat, 299
array_dims, 299
array_fill, 299
array_length, 299
array_lower, 299
array_ndims, 299
array_nrows configuration parameter, 585
array_position, 299
array_positions, 299
array_prepend, 299
array_remove, 299
array_replace, 299
array_to_json, 286
array_upper, 299
ASCII, 206
asin, 204
asind, 204
ASSERT
in PL/pgSQL, 1199
assertions
in PL/pgSQL, 1199
asynchronous commit, 742
AT TIME ZONE, 254
atan, 204
atan2, 204
atan2d, 204
autocommit
bulk-loading data, 458
psql, 1925
autovacuum
configuration parameters, 574
general information, 648
autovacuum configuration parameter, 574
autovacuum_analyze_scale_factor configuration parameter, 575
autovacuum_analyze_threshold configuration parameter, 574
autovacuum_freeze_max_age configuration parameter, 575
autovacuum_max_workers configuration parameter, 574
autovacuum_multixact_freeze_max_age configuration parameter, 575
autovacuum_naptime configuration parameter, 574
autovacuum_vacuum_cost_delay configuration parameter, 575
autovacuum_vacuum_cost_limit configuration parameter, 575
autovacuum_vacuum_scale_factor configuration parameter, 574
autovacuum_vacuum_threshold configuration parameter, 574
autovacuum_work_mem configuration parameter, 540
auto_explain.log_analyze configuration parameter, 3299
auto_explain.log_buffers configuration parameter, 3299
auto_explain.log_format configuration parameter, 3299
auto_explain.log_min_duration configuration parameter, 3299
auto_explain.log_nested_statements configuration parameter, 3300
auto_explain.log_timing configuration parameter, 3299
auto_explain.log_triggers configuration parameter, 3299
auto_explain.log_verbose configuration parameter, 3299
auto_explain.sample_rate configuration parameter, 3300
avals, 3355
average, 304
avg, 304

B-tree (see index)
backend_flush_after configuration parameter, 545
Background workers, 1329
backslash escapes, 34
backslash_quote configuration parameter, 585
backup, 339, 652
base type, 1028
BASE_BACKUP, 2110
BEGIN, 1462
BETWEEN, 199
BETWEEN SYMMETRIC, 200
BGWORKER_BACKEND_DATABASE_CONNECTION, 1330
BGWORKER_SHMEM_ACCESS, 1330
bgwriter_delay configuration parameter, 543
bgwriter_flush_after configuration parameter, 543
bgwriter_lru_maxpages configuration parameter, 543
bgwriter_lru_multiplier configuration parameter, 543
bigint, 38, 133
bigserial, 136
binary data, 140
functions, 219
binary string concatenation, 219
length, 220
bison, 475
Index

bit string
  constant, 37
  data type, 159
bit strings
  functions, 221
  bitmap scan, 375, 556
  bit_and, 304
  bit_length, 205
  bit_or, 304
BLOB (see large object)
  block_size configuration parameter, 587
  bloom, 3300
  bonjour configuration parameter, 535
  bonjour_name configuration parameter, 535
Boolean
  data type, 151
  operators (see operators, logical)
bool_and, 304
bool_or, 304
booting
  starting the server during, 506
box, 260
  box (data type), 155
BRIN (see index)
  brin_desummarize_range, 350
  brin_metapage_info, 3378
  brin_page_items, 3378
  brin_page_type, 3378
  brin_revmap_data, 3378
  brin_summarize_new_values, 350
  brin_summarize_range, 350
broadcast, 263
BSD Authentication, 612
btree_gin, 3304
btree_gist, 3304
btrim, 206, 220
bt_index_check, 3295
  bt_index_parent_check, 3296
bt_metap, 3376
  bt_page_items, 3377, 3377
  bt_page_stats, 3376
bytea, 140
bytea_output configuration parameter, 579
C
C, 770, 862
C++, 1070
CALL, 1464
canceling
  SQL command, 810
cardinality, 299
CASCADE
  with DROP, 99
  foreign key action, 66
Cascading Replication, 668
CASE, 295
  determination of result type, 368
case sensitivity

of SQL commands, 33
cast
  I/O conversion, 1495
cbrt, 202
  ceil, 202
  ceiling, 202
  center, 259
  Certificate, 611
  char, 138
  character, 138
  character set, 581, 588, 635
character string
  concatenation, 205
  constant, 34
  data types, 138
  length, 205
  character varying, 138
  char_length, 205
  check constraint, 60
  CHECK OPTION, 1633
  checkpoint, 743
CHECKPOINT, 1465
  checkpoint_completion_target configuration parameter, 549
  checkpoint_flush_after configuration parameter, 549
  checkpoint_timeout configuration parameter, 549
  checkpoint_warning configuration parameter, 550
  check_function_bodies configuration parameter, 577
  chr, 207
  cid, 194
  cidr, 157
  circle, 156, 261
citext, 3305
  client authentication, 594
  timeout during, 536
  client_encoding configuration parameter, 581
  client_min_messages configuration parameter, 566
  clock_timestamp, 245
CLOSE, 1466
class
  of databases (see database cluster)
CLUSTER, 1467
clusterdb, 1817
clustering, 668
  cluster_name configuration parameter, 572
cmax, 68
cmin, 67
COALESCE, 297
COLLATE, 50
collation, 629
  in PL/pgSQL, 1166
  in SQL functions, 1046
collation for, 328
column, 7, 58
  adding, 69
  removing, 69
  renaming, 71
  system column, 67
column data type
  changing, 70
column reference, 42
col_description, 334
column 
  about database objects, 334
  in SQL, 39
  COLUMN, 1469
  COMMIT, 1474
  COMMIT PREPARED, 1475
  column_delay configuration parameter, 549
  commit_siblings configuration parameter, 549
common table expression  (see WITH)
comparison
  composite type, 315
  operators, 198
  row constructor, 315
  subquery result row, 312
compiling
  libpq applications, 837
  composite type, 181, 1028
  comparison, 315
  constant, 183
  constructor, 52
  computed field, 187
  concat, 207
  concat_ws, 207
  concurrency, 426
conditional expression, 295
configuration
  of recovery
    of a standby server, 690
    of the server, 529
    of the server functions, 338
configure, 476
config_file configuration parameter, 533
conjunction, 198
connect by, 3435, 3442
connection service file, 831
conninfo, 777
constant, 34
constraint, 60
  adding, 69
  check, 60
  exclusion, 67
  foreign key, 64
  name, 60
  NOT NULL, 62
  primary key, 63
  removing, 70
  unique, 63
constraint exclusion, 97, 561
constraint_exclusion configuration parameter, 561
container type, 1028
CONTINUE
  in PL/pgSQL, 1183
  continuous archiving, 652
  in standby, 680
control file, 1100
convert, 207
convert_from, 207
convert_to, 208
COPY, 9, 1476
  with libpq, 813
corr, 306
correlation, 306
  in the query planner, 454
cos, 204
cosd, 204
cot, 204
cotd, 204
count, 304
covariance
  population, 306
  sample, 306
covar_pop, 306
covar_samp, 306
cpu_index_tuple_cost configuration parameter, 558
cpu_operator_cost configuration parameter, 558
cpu_tuple_cost configuration parameter, 558
CREATE ACCESS METHOD, 1486
CREATE AGGREGATE, 1487
CREATE CAST, 1495
CREATE COLLATION, 1499
CREATE CONVERSION, 1501
CREATE DATABASE, 621, 1503
CREATE DOMAIN, 1506
CREATE EVENT TRIGGER, 1509
CREATE EXTENSION, 1511
CREATE FOREIGN DATA WRAPPER, 1513
CREATE FOREIGN TABLE, 1515
CREATE FUNCTION, 1519
CREATE GROUP, 1527
CREATE INDEX, 1528
CREATE LANGUAGE, 1536
CREATE MATERIALIZED VIEW, 1539
CREATE OPERATOR, 1541
CREATE OPERATOR CLASS, 1544
CREATE OPERATOR FAMILY, 1547
CREATE POLICY, 1548
CREATE PROCEDURE, 1554
CREATE PUBLICATION, 1557
CREATE ROLE, 614, 1559
CREATE RULE, 1564
CREATE SCHEMA, 1567
CREATE SEQUENCE, 1570
CREATE SERVER, 1574
CREATE STATISTICS, 1576
CREATE SUBSCRIPTION, 1578
CREATE TABLE, 7, 1581
CREATE TABLE AS, 1601
CREATE TABLESPACE, 624, 1604
CREATE TEXT SEARCH CONFIGURATION, 1606
CREATE TEXT SEARCH DICTIONARY, 1607
CREATE TEXT SEARCH PARSER, 1609
CREATE TEXT SEARCH TEMPLATE, 1611
CREATE TRANSFORM, 1612
CREATE TRIGGER, 1614
CREATE TYPE, 1621
CREATE USER, 1630
CREATE USER MAPPING, 1631
CREATE VIEW, 1633
createdb, 3, 622, 1820
createuser, 614, 1823
CREATE_REPLICATION_SLOT, 2106
cross compilation, 483
cross join, 106
crosstab, 3436, 3438, 3439
crypt, 3385
cstring, 196
ctid, 68
CTID, 1136
CUBE, 117
cube (extension), 3308
cume_dist, 311
hypothetical, 310
current_catalog, 322
current_database, 322
current_date, 245
current_logfiles
    and the log_destination configuration parameter, 562
    and the pg_current_logfile function, 323
current_query, 322
current_role, 322
current_schema, 322
current_schemas, 322
current_setting, 338
current_time, 245
current_timestamp, 246
current_user, 322
currval, 293
cursor
    CLOSE, 1466
    DECLARE, 1639
    FETCH, 1706
    in PL/pgSQL, 1190
    MOVE, 1732
    showing the query plan, 1701
cursor_tuple_fraction configuration parameter, 561
custom scan provider
    handler for, 2179
Cygwin
    installation on, 493
D
data area (see database cluster)
data partitioning, 668
data type, 131
    base, 1028
    category, 360
    composite, 1028
    constant, 38
    container, 1028
    conversion, 359
    domain, 193
    enumerated (enum), 152
    internal organization, 1051
    numeric, 132
    polymorphic, 1029
    type cast, 49
    user-defined, 1077
database, 621
    creating, 3
    privilege to create, 615
database activity
    monitoring, 694
database cluster, 7, 504
data_checksums configuration parameter, 587
data_directory configuration parameter, 533
data_directory_mode configuration parameter, 587
date, 142, 144
    constants, 146
    current, 255
    output format, 146
    (see also formatting)
DateStyle configuration parameter, 580
date_part, 246, 249
date_trunc, 246, 253
dblink, 3313, 3319
dblink_build_sql_delete, 3341
dblink_build_sql_insert, 3339
dblink_build_sql_update, 3343
dblink_cancel_query, 3337
dblink_close, 3328
dblink_connect, 3314
dblink_connect_u, 3317
dblink_disconnect, 3318
dblink_error_message, 3331
dblink_exec, 3322
dblink_fetch, 3326
dblink_get_connections, 3330
dblink_get_notify, 3334
dblink_get_pkey, 3338
dblink_get_result, 3335
dblink_is_busy, 3333
dblink_open, 3324
dblink_send_query, 3332
db_user_namespace configuration parameter, 536
deadlock, 435
    timeout during, 584
deadlock_timeout configuration parameter, 584
DEALLOCATE, 1638
dearmorn, 3388
debug_assertions configuration parameter, 587
debug_deadlocks configuration parameter, 591
debug_pretty_print configuration parameter, 567
debug_print_parse configuration parameter, 567
debug_print_plan configuration parameter, 567
debug_print_rewritten configuration parameter, 567
decimal (see numeric)
DECLARE, 1639
decode, 208, 220
decode_bytea in PL/Perl, 1244
decrypt, 3392
decrypt_iv, 3392
default value, 59
  changing, 70
default_statistics_target configuration parameter, 560
default_tablespace configuration parameter, 576
default_text_search_config configuration parameter, 581
default_transaction_deferrable configuration parameter, 577
default_transaction_isolation configuration parameter, 577
default_transaction_read_only configuration parameter, 577
default_with_oids configuration parameter, 585
deferrable transaction
  setting, 1795
  setting default, 577
defined, 3356
degrees, 202
delay, 256
DELETE, 15, 103, 1642
  RETURNING, 103
delete, 3356
deleting, 103
dense_rank, 311
  hypothetical, 309
diameter, 259
dict_int, 3344
dict_xsyn, 3345
difference, 3350
digester, 3384
dirty read, 426
DISCARD, 1645
disjunction, 198
disk drive, 747
disk space, 643
disk usage, 738
DISTINCT, 10, 121
div, 202
dmetaphone, 3352
dmetaphone_alt, 3352
DO, 1646
document
  text search, 386
dollar quoting, 36
domain, 193
double precision, 135
DROP ACCESS METHOD, 1648
DROP AGGREGATE, 1649
DROP CAST, 1651
DROP COLLATION, 1652
DROP CONVERSION, 1653
DROP DATABASE, 624, 1654
DROP DOMAIN, 1655
DROP EVENT TRIGGER, 1656
DROP EXTENSION, 1657
DROP FOREIGN DATA WRAPPER, 1658
DROP FOREIGN TABLE, 1659
DROP FUNCTION, 1660
DROP GROUP, 1662
DROP INDEX, 1663
DROP LANGUAGE, 1665
DROP MATERIALIZED VIEW, 1666
DROP OPERATOR, 1667
DROP OPERATOR CLASS, 1669
DROP OPERATOR FAMILY, 1671
DROP OWNED, 1673
DROP POLICY, 1674
DROP PROCEDURE, 1675
DROP PUBLICATION, 1677
DROP ROLE, 614, 1678
DROP ROUTINE, 1679
DROP RULE, 1680
DROP SCHEMA, 1681
DROP SEQUENCE, 1682
DROP SERVER, 1683
DROP STATISTICS, 1684
DROP SUBSCRIPTION, 1685
DROP TABLE, 8, 1687
DROP TABLESPACE, 1688
DROP TEXT SEARCH CONFIGURATION, 1689
DROP TEXT SEARCH DICTIONARY, 1690
DROP TEXT SEARCH PARSER, 1691
DROP TEXT SEARCH TEMPLATE, 1692
DROP TRANSFORM, 1693
DROP TRIGGER, 1694
DROP TYPE, 1695
DROP USER, 1696
DROP USER MAPPING, 1697
DROP VIEW, 1698
dropdb, 624, 1827
donuser, 614, 1830
DROP_REPLICATION_SLOT, 2110
DTD, 164
DTrace, 483, 728
duplicate, 10
duplicates, 121
dynamic loading, 583, 1049
dynamic_library_path, 1050
dynamic_library_path configuration parameter, 583
dynamic_shared_memory_type configuration parameter, 541
E
each, 3356
earth, 3347
earthdistance, 3346
earth_box, 3347
earth_distance, 3347
ECPG, 862
ecpg, 1833
effective_cache_size configuration parameter, 559
effective_io_concurrency configuration parameter, 543
elog, 2139
   in PL/Perl, 1243
   in PL/Python, 1266
   in PL/Tcl, 1229
embedded SQL
   in C, 862
enabled role, 991
enable_bitmapscan configuration parameter, 556
enable_gathermerge configuration parameter, 556
enable_hashagg configuration parameter, 556
enable_hashjoin configuration parameter, 556
enable_indexonlyscan configuration parameter, 556
enable_indexscan configuration parameter, 556
enable_material configuration parameter, 556
enable_mergejoin configuration parameter, 556
enable_nestedloop configuration parameter, 556
enable_parallel_append configuration parameter, 556
enable_parallel_hash configuration parameter, 557
enable_partitionwise_aggregate configuration parameter, 557
enable_partitionwise_join configuration parameter, 557
enable_partition_pruning configuration parameter, 557
enable_seqscan configuration parameter, 557
enable_sort configuration parameter, 557
enable_tidscan configuration parameter, 557
encode, 208, 220
encode_array_constructor
   in PL/Perl, 1244
encode_array_literal
   in PL/Perl, 1244
encrypt, 3392
encryption, 522
for specific columns, 3384
encrypt_iv, 3392
END, 1699
enumerated types, 152
enum_first, 257
enum_last, 257
enum_range, 257
environment variable, 829
ephemeral named relation
   registering with SPI, 1301, 1303
   unregistering from SPI, 1302
ereport, 2138
error codes
   libpq, 795
   list of, 2292
error message, 787
escape string syntax, 34
escape_string_warning configuration parameter, 586
escaping strings
in libpq, 802
event log
   event log, 528
event trigger, 1120
   in C, 1125
   in PL/Tcl, 1231
event_source configuration parameter, 565
event_trigger, 196
every, 304
EXCEPT, 122
exceptions
   in PL/pgSQL, 1186
   in PL/Tcl, 1231
exclusion constraint, 67
EXECUTE, 1700
exist, 3356
EXISTS, 312
EXIT
   in PL/pgSQL, 1182
exit_on_error configuration parameter, 587
exp, 202
EXPLAIN, 441, 1701
expression
   order of evaluation, 54
   syntax, 41
extending SQL, 1028
extension, 1099
   externally maintained, 3468
external_pid_file configuration parameter, 533
extract, 246, 249
extra_float_digits configuration parameter, 581
F
failover, 668
false, 151
family, 263
fast path, 811
fdw_handler, 196
FETCH, 1706
field
   computed, 187
   field selection, 43
file system mount points, 505
file_fdw, 3348
FILTER, 44
first_value, 312
flex, 475
float4 (see real)
float8 (see double precision)
floating point, 135
floating-point
   display, 581
floor, 202
force_parallel_mode configuration parameter, 562
foreign data, 98
foreign data wrapper
   handler for, 2157
foreign key, 16, 64
Index

foreign table, 98
format, 208, 217
use in PL/pgSQL, 1172
formatting, 237
format_type, 328
Free Space Map, 2246
FreeBSD
IPC configuration, 511
shared library, 1058
start script, 507
from Collapse Limit configuration parameter, 561
FMS (see Free Space Map)
fsync configuration parameter, 546
full text search, 385
data types, 160
functions and operators, 160
full page writes configuration parameter, 547
function, 198
default values for arguments, 1039
in the FROM clause, 111
internal, 1049
invocation, 44
mixed notation, 57
named argument, 1031
named notation, 56
output parameter, 1037
polymorphic, 1029
positional notation, 56
RETURNS TABLE, 1044
type resolution in an invocation, 364
user-defined, 1030
in C, 1049
in SQL, 1030
variadic, 1038
with SETOF, 1040
functional dependency, 116
fuzzystrmatch, 3350

G
gc to sec, 3347
generate_series, 318
generate_subscripts, 319
genetic query optimization, 560
gen_random_bytes, 3392
gen_random_uuid, 3392
gen_salt, 3385
GEQO (see genetic query optimization)
geqo configuration parameter, 560
geqo_effort configuration parameter, 560
geqo_generations configuration parameter, 560
geqo_pool_size configuration parameter, 560
geqo_seed configuration parameter, 560
geqo_selection_bias configuration parameter, 560
geqo_threshold configuration parameter, 560
get_bit, 220
gt byte, 220
gt current ts config, 266
get raw page, 3374
GIN (see index)
gin_clean_pending_list, 350
gin_fuzzy_search_limit configuration parameter, 584
gin_leafpage_items, 3379
gin_metapage_info, 3379
gin_page_opaque_info, 3379
gin_pending_list_limit configuration parameter, 580
GiST (see index)
global data
in PL/Python, 1259
in PL/Tcl, 1226
GRANT, 71, 1710
GREATEST, 298
determination of result type, 368
Gregorian calendar, 2304
GROUP BY, 14, 115
grouping, 115
GROUPING, 310
GROUPING SETS, 117
gssapi, 604
GUID, 162

H
hash (see index)
hash_bitmap_info, 3381
hash_metapage_info, 3381
hash_page_items, 3380
hash_page_stats, 3380
hash_page_type, 3380
has any column privilege, 326
has column privilege, 326
has database privilege, 326
has foreign data wrapper privilege, 326
has function privilege, 326
has language privilege, 326
has schema privilege, 326
has sequence privilege, 326
has server privilege, 326
has tablespace privilege, 326
has table privilege, 326
has type privilege, 326
HAVING, 14, 116
hba file configuration parameter, 533
heap page_items, 3375
heap page item attrs, 3376
height, 259
hierarchical database, 7
high availability, 668
history
of PostgreSQL, xxxvi
hmac, 3384
host, 263
host name, 779
hostmask, 263
Hot Standby, 668
hot standby configuration parameter, 554
hot standby feedback configuration parameter, 554
isfinite, 246
isn, 3363
ISNULL, 200
isn_weak, 3365
isopen, 259
is_array_ref
  in PL/Perl, 1244
is_valid, 3365
J
JIT, 754
jit configuration parameter, 561
jit_above_cost configuration parameter, 559
jit_debugging_support configuration parameter, 592
jit_dump_bitcode configuration parameter, 592
jit_expressions configuration parameter, 592
jit_inline_above_cost configuration parameter, 559
jit_optimize_above_cost configuration parameter, 559
jit_profiling_support configuration parameter, 592
jit_provider configuration parameter, 583
jit_tuple_deforming configuration parameter, 592
join, 11, 106
  controlling the order, 456
  cross, 106
  left, 107
  natural, 108
  outer, 12, 107
  right, 107
  self, 13
join_collapse_limit configuration parameter, 561
JSON, 165
  functions and operators, 284
JSONB, 165
jsonb
  containment, 168
  existence, 168
  indexes on, 170
jsonb_agg, 304
jsonb_array_elements, 288
jsonb_array_elements_text, 288
jsonb_array_length, 288
jsonb_build_array, 286
jsonb_build_object, 286
jsonb_each, 288
jsonb_each_text, 288
jsonb_extract_path, 288
jsonb_extract_path_text, 288
jsonb_object, 286
jsonb_object_agg, 305
jsonb_object_keys, 288
jsonb_populate_record, 288
jsonb_populate_recordset, 288
jsonb_strip_nulls, 288
json_to_record, 288
json_to_recordset, 288
json_typeof, 288
json_agg, 304
json_array_elements, 288
json_array_elements_text, 288
json_array_length, 288
json_build_array, 286
json_build_object, 286
json_each, 288
json_each_text, 288
json_extract_path, 288
json_extract_path_text, 288
json_object, 286
json_object_agg, 305
json_object_keys, 288
json_populate_record, 288
json_populate_recordset, 288
json_strip_nulls, 288
json_to_record, 288
json_to_recordset, 288
json_typeof, 288
Julian date, 2304
Just-In-Time compilation (see JIT)
justify_days, 246
justify_hours, 246
justify_interval, 246
K
key word
  list of, 2306
  syntax of, 32
krb_caseins_users configuration parameter, 536
krb_server_keyfile configuration parameter, 536
L
label (see alias)
lag, 311
language_handler, 196
large object, 850
lastval, 293
last_value, 312
LATERAL
  in the FROM clause, 113
latitude, 3347
lea, 3371
lc_collate configuration parameter, 588
lc_cotype configuration parameter, 588
lc_messages configuration parameter, 581
lc_monetary configuration parameter, 581
lc_numeric configuration parameter, 581
lc_time configuration parameter, 581
LDAP, 480, 607
LDAP connection parameter lookup, 832
ldconfig, 489
lead, 311
LEAST, 298
  determination of result type, 368
left, 208
Index

ltree, 3368
ltree2text, 3371
ltrim, 209

M
MAC address (see macaddr)
MAC address (EUI-64 format) (see macaddr)
macaddr (data type), 158
macaddr8 (data type), 158
macaddr8_set7bit, 264
macOS
installation on, 495
IPC configuration, 513
shared library, 1058
magic block, 1049
maintenance, 642
maintenance_work_mem configuration parameter, 540
make, 474
make_date, 247
make_interval, 247
make_time, 247
make_timestamp, 247
make_timestampz, 247
make_valid, 3365
MANPATH, 489
masklen, 263
materialized view
implementation through rules, 1137
materialized views, 2073
max, 305
max_connections configuration parameter, 534
max_files_per_process configuration parameter, 541
max_function_args configuration parameter, 588
max_identifier_length configuration parameter, 588
max_index_keys configuration parameter, 588
max_locks_per_transaction configuration parameter, 584
max_logical_replication_workers configuration parameter, 555
max_parallel_maintenance_workers configuration parameter, 544
max_parallel_workers configuration parameter, 545
max_parallel_workers_per_gather configuration parameter, 544
max_pred_locks_per_page configuration parameter, 585
max_pred_locks_per_relation configuration parameter, 585
max_pred_locks_per_transaction configuration parameter, 584
max_prepared_transactions configuration parameter, 540
max_replication_slots configuration parameter, 551
max_stack_depth configuration parameter, 540
max_standby_archive_delay configuration parameter, 554
max_standby_streaming_delay configuration parameter, 554
max_sync_workers_per_subscription configuration parameter, 555
max_wal_senders configuration parameter, 551
max_wal_size configuration parameter, 551
max_worker_processes configuration parameter, 544
md5, 209, 220
MD5, 603
median, 46
memory context
in SPI, 1312
memory overcommit, 516
metaphone, 3352
min, 305
MinGW
installation on, 495
min_parallel_index_scan_size configuration parameter, 559
min_parallel_table_scan_size configuration parameter, 559
min_wal_size configuration parameter, 550
mod, 202
mode
statistical, 308
monitoring
database activity, 694
MOVE, 1732
moving-aggregate mode, 1072
Multiversion Concurrency Control, 426
MultiXactId, 648
MVCC, 426

N
name
qualified, 79
syntax of, 32
unqualified, 80
NaN (see not a number)
natural join, 108
negation, 198
NetBSD
IPC configuration, 512
shared library, 1058
start script, 507
netmask, 263
network, 263
data types, 156
Network Attached Storage (NAS) (see Network File Systems)
Network File Systems, 506
nextval, 293
NFS (see Network File Systems)
nlevel, 3371
non-durable, 461
nonblocking connection, 772, 805
nonrepeatable read, 426
normal_rand, 3435
NOT (operator), 198

3495
not a number
double precision, 135
numeric (data type), 134
NOT IN, 312, 315
not-null constraint, 62
notation
functions, 55
notice processing
in libpq, 822
notice processor, 822
notice receiver, 822
NOTIFY, 1734
in libpq, 812
NOTNULL, 200
now, 248
npoints, 259
nth_value, 312
ntile, 311
null value
with check constraints, 62
comparing, 200
default value, 59
in DISTINCT, 121
in libpq, 800
in PL/Perl, 1236
in PL/Python, 1254
with unique constraints, 63
NULLIF, 297
number
constant, 37
numeric, 38
numeric (data type), 133
numnode, 266, 400
num_nonnulls, 201
num_nulls, 201
NVL, 297

O
object identifier
data type, 194
object-oriented database, 7
obj_description, 334
octet_length, 205, 219
OFFSET, 123
OID
column, 67
in libpq, 802
oid, 194
oid2name, 3457
old_snapshot_threshold configuration parameter, 545
ON CONFLICT, 1719
ONLY, 106
OOM, 516
opaque, 196
OpenBSD
IPC configuration, 512
shared library, 1058
start script, 507

OpenSSL, 480
(see also SSL)
operator, 198
invocation, 43
logical, 198
precedence, 40
syntax, 38
type resolution in an invocation, 360
user-defined, 1081
operator class, 380, 1086
operator family, 380, 1094
operator_precedence_warning configuration parameter, 586
OR (operator), 198
Oracle
porting from PL/SQL to PL/pgSQL, 1214
ORDER BY, 10, 122
and locales, 628
ordered-set aggregate, 44
built-in, 308
ordering operator, 1096
ordinality, 320
outer join, 107
output function, 1077
OVER clause, 46
overcommit, 516
OVERLAPS, 248
overlay, 205, 219
overloading
functions, 1046
operators, 1082
owner, 71

P
pageinspect, 3374
page_checksum, 3375
page_header, 3375
palloc, 1057
PAM, 480, 611
parallel query, 462
parallel_leader_participation configuration parameter, 562
parallel_setup_cost configuration parameter, 559
parallel_tuple_cost configuration parameter, 559
parameter
syntax, 42
parenthesis, 42
parse_ident, 209
partition pruning, 96
partitioned table, 86
partitioning, 86
password, 615
authentication, 603
of the superuser, 505
password file, 831
passwordcheck, 3381
password_encryption configuration parameter, 536
path, 261
Index

for schemas, 575
PATH, 489
path (data type), 156
pattern matching, 222
patterns
  in psql and pg_dump, 1923
pclose, 259
peer, 607
percentile
  continuous, 308
  discrete, 308
percent_rank, 311
  hypothetical, 309
performance, 441
perl, 475
Perl, 1235
permission (see privilege)
pfree, 1057
PGAPPNAME, 830
pgbench, 1842
PGcancel, 810
PGCLIENTENCODING, 830
PGconn, 770
PGCONNECT_TIMEOUT, 830
pgcrypto, 3384
PGDATA, 504
PGDATABASE, 830
PGDATESTYLE, 830
PGEventProc, 825
PGEGO, 830
PGGSSLIB, 830
PGHOST, 829
PGHOSTADDR, 830
PGKRBSRVNAME, 830
PGLOCALEDIR, 831
PGOPTIONS, 830
PGPASSFILE, 830
PGPASSWORD, 830
PGPORT, 830
pgcrypto, 3384
pencrypt, 2010
pg_class, 2011
pg_client_encoding, 210
pg_collation, 2015
pg_collation_actual_version, 350
pg_collation_is_visible, 328
pg_column_size, 348
pg_config, 1858, 2067
  with ecpg, 921
  with libpq, 837
  with user-defined C functions, 1057
pg_conf_load_time, 323
pg_consume_ident, 316
pg_copy, 320
pg_convert_bytea, 3387
pg_convert, 3387
pg_convert_bytea, 3387
pg_conversion, 2019
PGDATA, 504
PGDATABASE, 830
PGDATESTYLE, 830
PGEventProc, 825
PGEGO, 830
PGGSSLIB, 830
PGHOST, 829
PGHOSTADDR, 830
PGKRBSRVNAME, 830
PGLOCALEDIR, 831
PGOPTIONS, 830
PGPASSFILE, 830
PGPASSWORD, 830
PGPORT, 830
PGREQUIREPEER, 830
PGREQUIRESSL, 830
PGresult, 793
pgrowlocks, 3397
pgstatinindex, 3405
pgstatashindex, 3406
pgstatindex, 3404
pgstatlock, 3403
pgstatlock_approx, 3406
PGSYSCONFDIR, 831
PGTARGETSESSIONATTRS, 830
PGTZ, 830
PGUSER, 830
pgxs, 1106
pg_advisory_lock, 354
pg_advisory_lock_shared, 354
pg_advisory_unlock, 354
pg_advisory_unlock_all, 354
pg_advisory_unlock_shared, 354
pg_advisory_xact_lock, 355
pg_advisory_xact_lock_shared, 355
pg_aggregate, 2000
pg_am, 2003
pg_amop, 2003
pg_amproc, 2004
pg_archivecleanup, 1949
pg_attrdef, 2009
pg_auth_members, 2010
pg_available_extensions, 2066
pg_available_extension_versions, 2067
pg_backend_pid, 322
pg_backup_start_time, 339
pg_basebackup, 1835
pg_blocking_pids, 323
pg_buffercache, 3382
pg_buffercache_pages, 3382
pg_cancel_backend, 339
pg_cast, 2010
pg_class, 2011
pg_client_encoding, 210
pg_collation, 2015
pg_collation_actual_version, 350
pg_collation_is_visible, 328
pg_column_size, 348
pg_config, 1858, 2067
  with ecpg, 921
  with libpq, 837
  with user-defined C functions, 1057
pg_conf_load_time, 323
pg_constraint, 2016
pg_copy, 320
pg_convert_bytea, 3387
pg_convert, 3387
pg_convert_bytea, 3387
pg_conversion, 2019
pg_conversion_is_visible, 328
pg_create_logical_replication_slot, 345
pg_create_physical_replication_slot, 344
pg_create_restore_point, 339
pg_ctl, 504, 506, 1952
pg_current logfile, 323
pg_current_wal_flush_lsn, 339
pg_current_wal_insert_lsn, 339
pg_current_wal_lsn, 339
pg_cursors, 2068
pg_database, 623, 2019
pg_database_size, 348
pg_db_role_setting, 2021
pg_ddl_command, 196
pg_default_acl, 2022
pg_description, 2024
pg_depend, 2022
pg_drop_replication_slot, 345
pg_dump, 1861
pg_dumpall, 1873
pg_event_trigger, 2025
pg_event_trigger_ddl_commands, 355
pg_event_trigger_dropped_objects, 356
pg_event_trigger_table_rewrite_oid, 358
pg_event_trigger_table_rewrite_reason, 358
pg_export_snapshot, 343
pg_extension, 2026
pg_extension_config_dump, 1103
pg_file_rename, 3294
pg_file_unlink, 3294
pg_file_write, 3294
pg_foreign_data_wrapper, 2026
pg_foreign_server, 2027
pg_foreign_table, 2028
pg_freespace, 3395
pg_freespacemap, 3395
pg_get_constraintdef, 328
pg_get_expr, 328
pg_get_functiondef, 328
pg_get_function_arguments, 328
pg_get_function_identity_arguments, 328
pg_get_function_result, 328
pg_get_indexdef, 328
pg_get_keywords, 328
pg_get_object_address, 333
pg_get_ruledef, 328
pg_get_serial_sequence, 328
pg_get_statisticsobjdef, 328
pg_get_triggerdef, 328
pg_get_userbyid, 328
pg_get_viewdef, 328
pg_group, 2069
pg_has_role, 326
pg_hba.conf, 594
pg_hba_file_rules, 2069
pg_ident.conf, 601
pg_identify_object, 333
pg_identify_object_as_address, 333
pg_import_system_collations, 350
pg_index, 2028
pg_indexexam_has_property, 328
pg_indexes, 2070
pg_indexes_size, 348
pg_index_column_has_property, 328
pg_index_has_property, 328
pg_inherits, 2031
pg_init_privs, 2031
pg_isready, 1879
pg_is_in_backup, 339
pg_is_in_recovery, 342
pg_is_other_temp_schema, 323
pg_is_wal_replay_paused, 343
pg_language, 2032
pg_largeobject, 2033
pg_largeobject_metadata, 2034
pg_last_committed_xact, 336
pg_last_wal_receive_lsn, 342
pg_last_wal_replay_lsn, 342
pg_last_xact_replay_timestamp, 342
pg_listening_channels, 323
pg_locks, 2071
pg_logdir_ls, 3294
pg_logical_emit_message, 347
pg_logical_slot_get_binary_changes, 345
pg_logical_slot_get_changes, 345
pg_logical_slot.peek_binary_changes, 346
pg_logical_slot.peek_changes, 345
pg_lsn, 196
pg_ls_dir, 352
pg_ls_logdir, 352
pg_ls_waldir, 352
pg_matviews, 2073
pg_my_temp_schema, 323
pg_namespace, 2034
pg_notification_queue_usage, 323
pg_notify, 1735
pg_opclass, 2034
pg_opclass_is_visible, 328
pg_operator, 2035
pg_operator_is_visible, 328
pg_opfamily, 2036
pg_opfamily_is_visible, 328
pg_options_to_table, 328
pg_policies, 2074
pg_policy, 2038
pg_postmaster_start_time, 324
pg_prepared_statements, 2075
pg_prepared_xacts, 2075
pg_prewarm, 3396
Index

pg_prewarm.autoprewarm configuration parameter, 3397
pg_prewarm.autoprewarm_interval configuration parameter, 3397
pg_proc, 2039
pg_publication, 2043
pg_publication_rel, 2044
pg_publication_table, 2076
pg_range, 2044
pg_read_binary_file, 352
pg_read_file, 352
pg_receivewal, 1881
pg_recvlogical, 1885
pg_relation_filenode, 349
pg_relation_filepath, 349
pg_relation_size, 348
pg_reload_conf, 339
pg_relpages, 3406
pg_replication_origin, 2045
pg_replication_origin_advance, 347
pg_replication_origin_create, 346
pg_replication_origin_drop, 346
pg_replication_origin_oid, 346
pg_replication_origin_progress, 347
pg_replication_origin_session_is_setup, 346
pg_replication_origin_session_progress, 346
pg_replication_origin_session_reset, 346
pg_replication_origin_status, 2076
pg_replication_origin_xact_reset, 347
pg_replication_origin_xact_setup, 347
pg_replication_slots, 2077
pg_replication_slot_advance, 346
pg_resetwal, 1958
pg_restore, 1889
pg_rewind, 1961
pg_rewrite, 2045
pg_roles, 2078
pg_rotate_logfile, 339
pg_rules, 2079
pg_safe_snapshot_blocking_pids, 324
pg_seclabel, 2046
pg_seclabels, 2080
pg_sequence, 2046
pg_sequences, 2080
pg_service.conf, 831
pg_settings, 2081
pg_shadow, 2083
pg_shdepend, 2047
pg_shdescription, 2048
pg_shseclabel, 2049
pg_size_bytes, 348
pg_size_pretty, 348
pg_sleep, 256
pg_sleep_for, 256
pg_sleep_until, 256
pg_standby, 3463
pg_start_backup, 339
pg_statio_all_indexes, 698
pg_statio_all_sequences, 698
pg_statio_all_tables, 698
pg_statio_sys_indexes, 698
pg_statio_sys_sequences, 698
pg_statio_sys_tables, 698
pg_statio_user_indexes, 698
pg_statio_user_sequences, 698
pg_statio_user_tables, 698
pg_statistic, 453, 2049
pg_statistics_obj_is_visible, 328
pg_statistic_ext, 454, 2051
pg_stats, 453, 2084
pg_stat_activity, 696
pg_stat_all_indexes, 697
pg_stat_all_tables, 697
pg_stat_archiver, 697
pg_stat_bgwriter, 697
pg_stat_clear_snapshot, 724
pg_stat_database, 697
pg_stat_database_conflicts, 697
pg_stat_file, 353
pg_stat_get_activity, 724
pg_stat_get_snapshot_timestamp, 724
pg_stat_progress_vacuum, 697
pg_stat_replication, 696
pg_stat_reset, 724
pg_stat_reset_shared, 724
pg_stat_reset_single_function_counters, 725
pg_stat_reset_single_table_counters, 725
pg_stat_ssl, 697
pg_stat_statements, 3398
  function, 3401
pg_stat_statements_reset, 3401
pg_stat_subscription, 697
pg_stat_sys_indexes, 697
pg_stat_sys_tables, 697
pg_stat_user_functions, 698
pg_stat_user_indexes, 697
pg_stat_user_tables, 697
pg_stat_wal_receiver, 696
pg_stat_xact_all_tables, 697
pg_stat_xact_sys_tables, 697
pg_stat_xact_user_functions, 698
pg_stat_xact_user_tables, 697
pg_stop_backup, 339
pg_subscription, 2052
pg_subscription_rel, 2053
pg_switch_wal, 339
pg_temp, 576
pg_tempf, 576
pg_tempf_size, 348
pg_table_is_visible, 328
pg_table_size, 348
pg_table_is_visible, 328
pg_table_size, 348
pg_table_size, 348
securing functions, 1525
Index

pg_terminate_backend, 339
pg_test_fsync, 1964
pg_test_timing, 1965
pg_timezone_abbrevs, 2087
pg_timezone_names, 2087
pg_total_relation_size, 348
pg_transform, 2054
pg_trgm, 3407
pg_trgm.similarity_threshold configuration parameter, 3410
pg_trgm.word_similarity_threshold configuration parameter, 3410
pg_trigger, 2054
pg_try_advisory_lock, 354
pg_try_advisory_lock_shared, 354
pg_try_advisory_xact_lock, 355
pg_try_advisory_xact_lock_shared, 355
pg_ts_config, 2056
pg_ts_config_is_visible, 328
pg_ts_config_map, 2056
pg_ts_dict, 2057
pg_ts_dict_is_visible, 328
pg_ts_parser, 2057
pg_ts_parser_is_visible, 328
pg_ts_template, 2058
pg_ts_template_is_visible, 328
pg_type, 2058
pg_typeof, 328
pg_type_is_visible, 328
pg_upgrade, 1969
pg_user, 2088
pg_user_mapping, 2065
pg_user_mappings, 2088
pg_verify_checksums, 1977
pg_views, 2089
pg_visibility, 3413
pg_waldump, 1978
pg_walfile_name, 339
pg_walfile_name_offset, 339
pg_wal_lsn_diff, 339
pg_wal_replay_pause, 343
pg_wal_replay_resume, 343
pg_xact_commit_timestamp, 336
phantom read, 426
phraseto_tsquery, 266, 393
pi, 202
PIC, 1057
PID
determining PID of server process
   in libpq, 787
PITR, 652
PITR standby, 668
pkg-config, 480
   with ecpg, 921
   with libpq, 838
PL/Perl, 1235
PL/PerlU, 1246
PL/pgSQL, 1159
PL/Python, 1251
PL/SQL (Oracle)
   porting to PL/pgSQL, 1214
PL/Tcl, 1224
plainto_tsquery, 266, 393
plperl.on_init configuration parameter, 1248
plperl.on_plperl_init configuration parameter, 1249
plperl.on_plperlu_init configuration parameter, 1249
plperl.use_strict configuration parameter, 1249
plpgsql.check asserts configuration parameter, 1199
plpgsql.variable_conflict configuration parameter, 1209
pltcl.start_proc configuration parameter, 1234
pltclu.start_proc configuration parameter, 1234
point, 154, 261
point-in-time recovery, 652
policy, 72
polygon, 156, 261
polymorphic function, 1029
polymorphic type, 1029
popen, 259
populate_record, 3356
port, 779
port configuration parameter, 534
position, 205, 219
POSTGRES, xxxvii
postgres, 3, 506, 622, 1980
postgres user, 504
Postgres95, xxxvii
postgresql.auto.conf, 530
postgresql.conf, 529
postgres_fdw, 3414
postmaster, 1987
post_auth_delay configuration parameter, 589
power, 202
PQbackendPID, 787
PQbinaryTuples, 799
   with COPY, 813
PQcancel, 810
PQclear, 797
PQclientEncoding, 817
PQcmdStatus, 801
PQcmdTuples, 801
PQconninfo, 775
PQconninfoFree, 819
PQconninfoParse, 775
PQconnmultiplex, 808
PQconnectdb, 771
PQconnectdbParams, 771
PQconsumeInput, 808
PQcopyResult, 820
PQdb, 784
PQdescribeport, 793
PQdescribePrepared, 792
PQencryptPassword, 820
PQencryptPasswordConn, 819
PQendcopy, 816
PQerrorMessage, 787
PQescapeBytea, 804
PQescapeByteaConn, 804
PQescapeIdentifier, 803
PQescapeLiteral, 802
PQescapeString, 804
PQescapeStringConn, 803
PQexec, 789
PQexecParams, 790
PQexecPrepared, 792
PQfformat, 799
PQfinish, 775
PQfireResultCreateEvents, 820
PQflush, 809
PQfmod, 799
PQfn, 811
PQfname, 798
PQfnumber, 798
PQfreeCancel, 810
PQfreemem, 819
PQfsize, 799
PQhtable, 798
PQhtablecol, 798
PQftype, 799
PQgetCancel, 810
PQgetCopyData, 814
PQgetisnull, 800
PQgetlength, 800
PQgetline, 815
PQgetlineAsync, 815
PQgetResult, 807
PQgetssl, 789
PQgetvalue, 800
PQhost, 784
PQinitOpenSSL, 836
PQinitSSL, 836
PQinstanceData, 826
PQisBusy, 808
PQisnonblocking, 809
PQisthreadsafe, 836
PQlibVersion, 821
PQmakeEmptyPGresult, 820
PQnfields, 797
PQnparams, 800
PQoidStatus, 802
PQoidValue, 802
PQoptions, 785
PQparameterStatus, 785
PQpass, 784
PQping, 777
PQpingParams, 776
PQport, 784
PQprepare, 791
PQprint, 801
PQprotocolVersion, 786
PQputCopyData, 814
PQputCopyEnd, 814
PQputnbytes, 816
PQregisterEventProc, 826
PQrequestEventProc, 826
Preset, 775
PresetCancel, 811
PresetPoll, 776
PresetStart, 776
PqresStatus, 794
PQresultAlloc, 821
PQresultErrorField, 795
PQresultErrorMessage, 794
PQresultInstanceData, 826
PQresultSetInstanceData, 826
PQresultVerboseErrorMessage, 794
PQsendDescribePortal, 807
PQsendDescribePrepared, 807
PQsendPrepare, 806
PQsendQuery, 805
PQsendQueryParams, 806
PQsendQueryPrepared, 806
PQserverVersion, 786
PQsetClientEncoding, 817
PQsetdb, 772
PQsetdbLogin, 771
PQsetErrorContextVisibility, 818
PQsetErrorVerbosity, 817
PQsetInstanceData, 826
PQsetnonblocking, 808
PQsetNoticeProcessor, 822
PQsetNoticeReceiver, 822
PQsetResultAttrs, 821
PQsetSingleRowMode, 809
PQsocket, 787
PQsslAttribute, 788
PQsslAttributeNames, 788
PQsslInUse, 788
PQsslStruct, 788
PQstatus, 785
PQtrace, 818
PQtransactionStatus, 785
PQtty, 785
PQunescapeBytea, 805
PQuntrace, 818
PQuser, 784
predicate locking, 430
PREPARE, 1737
PREPARE TRANSACTION, 1740

3501
prepared statements
  creating, 1737
  executing, 1700
  removing, 1638
  showing the query plan, 1701
preparing a query
  in PL/pgSQL, 1210
  in PL/Python, 1261
  in PL/Tcl, 1228
pre_auth_delay configuration parameter, 589
primary key, 63
primary_conninfo recovery parameter, 692
primary_slot_name recovery parameter, 693
privilege, 71
  querying, 324
  with rules, 1151
  for schemas, 81
  with views, 1151
procedural language, 1156
  externally maintained, 3468
  handler for, 2154
procedure
  user-defined, 1030
protocol
  frontend-backend, 2090
  ps
    to monitor activity, 694
psql, 5, 1898
Python, 1251
Q
  qualified name, 79
query, 9, 105
query plan, 441
query tree, 1129
querytree, 266, 400
  quotation marks
    and identifiers, 33
    escaping, 34
quote_all_identifiers configuration parameter, 586
quote_ident, 210
  in PL/Perl, 1244
  use in PL/pgSQL, 1172
quote_literal, 210
  in PL/Perl, 1243
  use in PL/pgSQL, 1172
quote_nullable, 210
  in PL/Perl, 1243
  use in PL/pgSQL, 1172
random_page_cost configuration parameter, 558
random, 204
range table, 1129
range type, 188
  exclude, 192
  indexes on, 192
rank, 311
  hypothetical, 309
read committed, 427
read-only transaction
  setting, 1795
  setting default, 577
realtime, 474
real, 135
REASSIGN OWNED, 1742
record, 196
recovery.conf, 690
recovery_end_command recovery parameter, 691
recovery_min_apply_delay recovery parameter, 693
recovery_target recovery parameter, 691
recovery_target_action recovery parameter, 692
recovery_target_inclusive recovery parameter, 691
recovery_target_lsn recovery parameter, 691
recovery_target_name recovery parameter, 691
recovery_target_timeline recovery parameter, 692
recovery_target_xid recovery parameter, 691
rectangle, 155
RECURSIVE
  in common table expressions, 125
  in views, 1633
referential integrity, 16, 64
REFRESH MATERIALIZED VIEW, 1743
regclass, 194
regconfig, 194
regdictionary, 194
regexp_match, 211, 224
regexp_matches, 211, 224
regexp_replace, 211, 224
regexp_split_to_array, 211, 224
regexp_split_to_table, 211, 224
regexp, 194
regoperator, 194
regprocedure, 194
regression intercept, 307
regression slope, 307
regression test, 487
regression tests, 757
regr_avgx, 306
regr_avgy, 306
regr_count, 306
regr_intercept, 307
regr_r2, 307
regr_slope, 307
regr_sxx, 307
regr_sxy, 307
regr_syy, 307
regtype, 194
regular expression, 223, 224
   (see also pattern matching)
regular expressions
   and locales, 628
reindex, 649
REINDEX, 1745
reindexdb, 1937
relation, 7
relational database, 7
RELEASE SAVEPOINT, 1748
repeat, 211
repeatable read, 429
replace, 212
replication, 668
Replication Origins, 1342
Replication Progress Tracking, 1342
replication slot
   logical replication, 1335
   streaming replication, 675
reporting errors
   in PL/pgSQL, 1197
RESET, 1750
restartpoint, 745
restart_after_crash configuration parameter, 587
restore_command recovery parameter, 690
RESTRICT
   with DROP, 99
   foreign key action, 66
RETURN NEXT
   in PL/pgSQL, 1176
RETURN QUERY
   in PL/pgSQL, 1176
RETURNING, 103
RETURNING INTO
   in PL/pgSQL, 1169
reverse, 212
REVOKE, 71, 1751
right, 212
right join, 107
role, 614, 618
   applicable, 970
   enabled, 991
   membership in, 616
   privilege to create, 615
   privilege to initiate replication, 615
ROLLBACK, 1755
rollback
   psql, 1927
ROLLBACK PREPARED, 1756
ROLLBACK TO SAVEPOINT, 1757
ROLLUP, 117
round, 203
routine, 1030
routine maintenance, 642
row, 7, 58
ROW, 52
row estimation
   multivariate, 2263
   planner, 2258
row type, 181
   constructor, 52
row-level security, 72
row-wise comparison, 315
row_number, 311
row_security configuration parameter, 576
row_security_active, 326
row_to_json, 286
rpad, 212
rtrim, 212
rule, 1129
   and materialized views, 1137
   and views, 1130
   for DELETE, 1140
   for INSERT, 1140
   for SELECT, 1131
   compared with triggers, 1153
   for UPDATE, 1140
S
SAVEPOINT, 1759
savepoints
   defining, 1759
   releasing, 1748
   rolling back, 1757
scalar (see expression)
   scale, 203
schema, 78, 621
   creating, 79
   current, 80, 322
   public, 79
   removing, 79
SCRAM, 603
search path, 80
   current, 322
   object visibility, 327
search_path configuration parameter, 80, 575
   use in securing functions, 1525
SECURITY LABEL, 1761
sec_to_gc, 3347
seg, 3420
segment_size configuration parameter, 588
SELECT, 9, 105, 1764
   determination of result type, 370
   select list, 120
SELECT INTO, 1785
   in PL/pgSQL, 1169
semaphores, 509
sepgsql, 3423
sepgsql.debug_audit configuration parameter, 3426
sepgsql.permissive configuration parameter, 3426
sequence, 293
   and serial type, 136
sequential scan, 557
seq_page_cost configuration parameter, 558
serial, 136
serial2, 136
serial4, 136
serial8, 136
serializable, 430
Serializable Snapshot Isolation, 426
serialization anomaly, 427, 430
server log, 562
server spoofing, 522
server_encoding configuration parameter, 588
server_version configuration parameter, 588
server_version_num configuration parameter, 588
session_preload_libraries configuration parameter, 582
session_replication_role configuration parameter, 578
session_user, 322
SET, 338, 1787
SET CONSTRAINTS, 1790
set difference, 122
set intersection, 122
set operation, 122
set returning functions
functions, 318
SET ROLE, 1791
SET SESSION AUTHORIZATION, 1793
SET TRANSACTION, 1795
set union, 122
SET XML OPTION, 580
setseed, 204
setval, 293
setweight, 266, 399
setweight for specific lexeme(s), 266
set_bit, 220
set_byte, 221
set_config, 338
set_limit, 3409
set_masklen, 263
sha224, 221
sha256, 221
sha384, 221
sha512, 221
shared library, 488, 1057
shared memory, 509
shared_buffers configuration parameter, 539
shared_preload_libraries, 1069
shared_preload_libraries configuration parameter, 583
shobj_description, 334
SHOW, 338, 1798, 2106
show_limit, 3408
show_trgm, 3408
shutdown, 518
SIGHUP, 530, 599, 602
SIGINT, 518
sign, 203
signal
backend processes, 339
significant digits, 581
SIGQUIT, 518
SIGTERM, 518
SIMILAR TO, 223
similarity, 3408
sin, 204
sind, 204
single-user mode, 1983
skeys, 3355
sleep, 256
slice, 3356
sliced bread (see TOAST)
smallint, 133
smallserial, 136
Solaris
installation on, 496
IPC configuration, 514
shared library, 1059
start script, 508
SOME, 305, 312, 315
sort, 3361
sorting, 122
sort_asc, 3361
sort_desc, 3361
soundex, 3350
SPI-GiST (see index)
SPI, 1269
examples, 3431
spi_commit
in PL/Perl, 1243
SPI_commit, 1323
SPI_connect, 1270
SPI_connect_ext, 1270
SPI_cursor_close, 1298
SPI_cursor_fetch, 1294
SPI_cursor_find, 1293
SPI_cursor_move, 1295
SPI_cursor_open, 1289
SPI_cursor_open_with_args, 1290
SPI_cursor_open_with_paramlist, 1292
SPI_exec, 1275
SPI_execp, 1288
SPI_execute, 1272
SPI_execute_plan, 1285
SPI_execute_plan_with_paramlist, 1287
SPI_execute_with_args, 1276
spi_exec_prepared
in PL/Perl, 1241
spi_exec_query
in PL/Perl, 1239
spi_fetchrow
in PL/Perl, 1240
SPI_finish, 1271
SPI_fnumber, 1305
SPI_freeplan
in PL/Perl, 1241
SPI_freeplan, 1322
SPI_freetuple, 1320
substr, 212
substring, 205, 219, 223, 224
subtransactions
in PL/Tcl, 1232
sum, 305
superuser, 5, 615
superuser_reserved_connections configuration parameter, 534
support functions
in_range, 2205
suppress_redundant_updates_trigger, 355
svals, 3355
synchronize_seqscans configuration parameter, 586
synchronous commit, 742
Synchronous Replication, 668
synchronous_commit configuration parameter, 546
synchronous_standby_names configuration parameter, 552
syntax
SQL, 32
syslog_facility configuration parameter, 565
syslog_ident configuration parameter, 565
date format, 565
date format, 565
date format, 565
syslog_sequence_numbers configuration parameter, 565
date format, 565
date format, 565
syslog_split_messages configuration parameter, 565
date format, 565
date format, 565
system catalog
date format, 81
date format, 81
date format, 81
date format, 81
RemoveIPC, 514
T
table, 7, 58
creating, 58
inheritance, 83
modifying, 68
partitioning, 86
removing, 59
renaming, 71
TABLE command, 1764
table expression, 105
table function, 111
XMLTABLE, 278
table sampling method, 2176
tablefunc, 3435
tableoid, 67
TABLESAMPLE method, 2176
tablespace, 624
default, 576
temporary, 577
tan, 204
tand, 204
target list, 1130
Tcl, 1224
tcn, 3444
tcp_keepalives_count configuration parameter, 535
tcp_keepalives_idle configuration parameter, 535
tcp_keepalives_interval configuration parameter, 535
template0, 622
template1, 622, 622
temp_buffers configuration parameter, 539
temp_file_limit configuration parameter, 541
temp_tablespaces configuration parameter, 577
test, 757
test_decoding, 3446
text, 138, 263
text search, 385
data types, 160
functions and operators, 160
indexes, 421
text2ltree, 3371
threads
with libpq, 836
tid, 194
time, 142, 144
constants, 146
current, 255
output format, 146
(time span, 142
time with time zone, 142, 144
time without time zone, 142, 144
time zone, 147, 580
conversion, 254
input abbreviations, 2303
time zone data, 482
time zone names, 581	
timelines, 652
TIMELINE_HISTORY, 2106
timeofday, 248
timeout
client authentication, 536
deadlock, 584
timestamp, 142, 145
timestamp with time zone, 142, 145
timestamp without time zone, 142, 145
timestamptz, 142
TimeZone configuration parameter, 580
timezone_abbreviations configuration parameter, 581
TOAST, 2243
and user-defined types, 1080
per-column storage settings, 1429
versus large objects, 850
token, 32
to_ascii, 213
to_char, 237
and locales, 628
to_date, 237
to_hex, 213
to_json, 286
to_jsonb, 286
to_number, 238
to_regclass, 328
to_regnamespace, 328
to_regoperator, 328
to_regproc, 328

3506
to_regprocedure, 328
to_regrole, 328
to_regtype, 328
to_timestamp, 238, 248
to_tsquery, 266, 392
to_tsvector, 266, 391
trace_locks configuration parameter, 590
trace_lock_oidmin configuration parameter, 590
trace_lwlocks configuration parameter, 590
trace_notify configuration parameter, 589
trace_recovery_messages configuration parameter, 589
trace_sort configuration parameter, 590
trace_userlocks configuration parameter, 590
track_activities configuration parameter, 573
track_activity_query_size configuration parameter, 573
track_commit_timestamp configuration parameter, 552
track_counts configuration parameter, 573
track_functions configuration parameter, 573
track_io_timing configuration parameter, 573
transaction, 17
transaction ID
wraparound, 645
transaction isolation, 426
transaction isolation level, 427
read committed, 427
repeatable read, 429
serializable, 430
setting, 1795
setting default, 577
transaction log (see WAL)
transaction_timestamp, 248
transform_null_equals configuration parameter, 586
transition tables, 1614
(see also ephemeral named relation)
implementation in PLs, 1303
referencing from C trigger, 1113
translate, 213
transparent huge pages, 539
trigger, 196, 1110
arguments for trigger functions, 1112
for updating a derived tvector column, 401
in C, 1113
in PL/pgSQL, 1199
in PL/Python, 1259
in PL/Tcl, 1229
compared with rules, 1153
triggered_change_notification, 3444
trigger_file recovery parameter, 693
true, 151
trunc, 203, 264, 264
TRUNCATE, 1801
trusted
PL/Perl, 1245
tsm_handler, 196
tsm_system_rows, 3446
tsm_system_time, 3446
tquery (data type), 161
tquery_phrase, 269, 399
tvector (data type), 160
tvector_concatenation, 398
tvector_to_array, 269
tvector_update_trigger, 269
tvector_update_trigger_column, 269
ts_debug, 270, 416
ts_delete, 267
ts_filter, 268
ts_headline, 268, 397
ts_lexize, 270, 420
ts_parse, 270, 419
ts_rank, 268, 395
ts_rank_cd, 268, 395
ts_rewrite, 268, 400
ts_stat, 271, 403
ts_token_type, 270, 419
tuple_data_split, 3376
txid_current, 335
txid_current_if_assigned, 335
txid_current_snapshot, 335
txid_snapshot_xip, 335
txid_snapshot_xmax, 335
txid_snapshot_xmin, 335
txid_status, 335
txid_visible_in_snapshot, 335
type (see data type)
type cast, 38, 49

U
UESCAPE, 33, 36
unaccent, 3447, 3449
Unicode escape
in identifiers, 33
in string constants, 35
UNION, 122
determination of result type, 368
uniq, 3361
unique constraint, 63
Unix domain socket, 779
unix_socket_directories configuration parameter, 534
unix_socket_group configuration parameter, 534
unix_socket_permissions configuration parameter, 534
unknown, 196
UNLISTEN, 1803
unnest, 299
for tvector, 269
unqualified name, 80
updatable views, 1635
UPDATE, 15, 102, 1805
RETURNING, 103
update_process_title configuration parameter, 572
upgrading, 102
upgrading, 519
upper, 206, 303
and locales, 628
upper_inc, 303
upper_inf, 303
UPSERT, 1719
URI, 777
user, 322, 614
current, 322
user mapping, 98
User name maps, 601
UUID, 162, 481
uuid-oss, 3449
uuid_generate_v1, 3449
uuid_generate_v1mc, 3449
uuid_generate_v3, 3449
V
vacuum, 642
VACUUM, 1810
vacuumdb, 1940
vacuumlo, 3461
vacuum_cleanup_index_scale_factor configuration parameter, 579
vacuum_cost_delay configuration parameter, 542
vacuum_cost_limit configuration parameter, 542
vacuum_cost_page_dirty configuration parameter, 542
vacuum_cost_page_hit configuration parameter, 542
vacuum_cost_page_miss configuration parameter, 542
vacuum_defer_cleanup_age configuration parameter, 553
vacuum_freeze_min_age configuration parameter, 579
vacuum_freeze_table_age configuration parameter, 578
vacuum_multixact_freeze_min_age configuration parameter, 579
vacuum_multixact_freeze_table_age configuration parameter, 579
value expression, 41
VALUES, 124, 1813
WHERE, 114
WHERE to log, 562
WHILE
in PL/pgSQL, 1183
width, 259
width_bucket, 203
window function, 19
built-in, 310
invocation, 46
order of execution, 120
WITH
in SELECT, 125, 1764
WITH CHECK OPTION, 1633
WITHIN GROUP, 44
xid, 194
xmax, 67
xmin, 67
XML, 163
XML export, 281
XML option, 164, 580
xml2, 3451
xmlagg, 275, 305
xmlelement, 272
xmlconcat, 271
xmlcomment, 271
xmlconcat, 271
xmlcomment, 271
Index

XMLEXISTS, 276
xmlforest, 273
xmloption configuration parameter, 580
xmlparse, 163
xmlpi, 274
xmlroot, 274
xmlserialize, 164
xmltable, 278
xml_is_well_formed, 276
xml_is_well_formed_content, 276
xml_is_well_formed_document, 276
XPath, 277
xpath_exists, 278
xpath_table, 3452
xslt_process, 3454

Y
yacc, 475

Z
zero_damaged_pages configuration parameter, 591
zlib, 475, 483